Power and Bandwidth Minimization for Demand-Aware GEO Satellite Systems

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Abstract—Smart radio resource allocation combined with the recent advances of digital payloads will allow to control the transmit power and bandwidth of the satellites depending on the demand and the channel conditions of users. The system flexibility is important not only to handle divergent demand requirements but also to efficiently utilize the limited and expensive satellite resources. In this paper, we propose a demand-aware smart radio resource allocation technique, where the transmit power and the bandwidth of the GEO satellite are minimized while satisfying the user demand. The formulated optimization problem is non-convex mixed-integer nonlinear program which is difficult to solve. Hence, we apply a quadratic transform to solve the problem iteratively. The numerical results showed that the proposed scheme outperforms the benchmark schemes in terms of bandwidth utilization while accurately providing capacity-on-demand.

Index Terms—GEO satellite, Smart radio resource allocation, Demand satisfaction, Bandwidth utilization, Quadratic transform

I. INTRODUCTION

The future satellites are expected to provide a high throughput to serve the increasing connectivity and traffic demands. Furthermore, the deployment of advanced digital payloads on these satellites will improve the control and the flexibility of resource allocation [1]. SES-17 [2] and EUTELSAT QUANTUM [3] are typical examples of such payloads. To fully utilize the scarce and expensive satellite resources as well as to support throughput-demanding services, smart radio resource allocation on the digital payloads is essential [4]. For this, a thorough research has been conducted to adaptively control the power and/or the bandwidth of the system with the aim of demand satisfaction. The resource allocation techniques developed so far can be divided into three categories:

1) Power Allocation: The system power is optimized based on the beams’ requested demand while considering a uniform bandwidth allocation. In [5]–[7], a power allocation was studied to minimize the difference between the offered capacity and the requested demand. However, the bandwidth allocation is considered to be fixed where a conventional frequency reuse pattern is applied at the cost of spectrum utilization. In [8], a power optimization technique based on channel condition is considered, while in [9], power allocation with imperfect channel state information (CSI) is proposed. A distinct drawback of these works is that inter-beam interference has been neglected in these works. On the other hand, power allocation with the full frequency reuse scheme is considered in [10].

2) Bandwidth allocation: In this case, the system bandwidth is optimized assuming a fixed power to satisfy the requested demand. In this context, bandwidth assignment with the assumption of time sharing for Gaussian broadcast channels has been addressed in [11], [12]. In these works, the interference among the beams has not been considered again.

3) Joint power and bandwidth allocation: Both power and bandwidth are optimized to satisfy the requested demand. This joint allocation has been introduced in [13], [14], where orthogonal carriers are assigned within a cluster of beams. Moreover, a binary power allocation is considered for each carrier. Similarly, to minimize the unmet system capacity, an evolutionary algorithms have been proposed in [15], [16]. Furthermore, in [17], a joint carrier and power assignment is considered for demand matching. However, the optimization solution is infeasible when the system offered capacity is below the required demand. Hence, lower demand satisfaction is occurred.

In this paper, we consider power and bandwidth minimization for demand-aware GEO Satellite Systems. Hence, carrier and power are allocated wherever demand is needed. Furthermore, we consider an interference-limited GEO satellite system with a large number of beams sharing the spectrum resources. The formulated optimization problem is non-convex mixed-integer nonlinear program which is difficult to solve. Therefore, we apply a quadratic transform to solve the problem iteratively. The extensive numerical results show that the demand matching error and bandwidth utilization of the proposed strategy is substantially lower than that obtained with the benchmark schemes.

The rest of the paper is organized as follows. Section II describes the system model. The problem formulation and the proposed solution are provided in Section III and Section IV, respectively. The numerical results are presented in Section V and Section VI concludes the paper.
II. SYSTEM MODEL

Consider a broadband geostationary (GEO) satellite system with \(N\) beams. The total bandwidth \(B_{\text{tot}}\) is subdivided into \(K\) subcarrier bandwidths of \(B_{\text{sc}}\) i.e. \(B_{\text{sc}} = B_{\text{tot}}/K\). Furthermore, we define a carrier allocation vector for carrier \(k\) as \(x_k = [x_k[1], x_k[2], \ldots, x_k[N]]^T\), where \(x_k[i]\) is the binary-valued variable and \(x_k[i] = 1\) means that the carrier \(k\) is allocated to the \(i\)th beam. Similarly, we define the power allocation vector for carrier \(k\) as \(p_k = [p_k[1], p_k[2], \ldots, p_k[N]]^T\), where \(p_k[i]\) is the power required by the \(i\)th beam to transmit in the \(k\)th carrier. Assuming single user per beam, the signal-to-interference-plus-noise ratio (SINR) of the \(i\)th beam over \(k\)th carrier is given by

\[
\gamma_{i,k}(p_k, x_k) = \frac{g_{i,k}[1]p_k[i]x_k[i]}{\sum_{j=1, j \neq i}^{N} g_{i,k}[j]p_k[j]x_k[j] + \sigma^2}, \tag{1}
\]

where \(g_{i,k}[i]\) the channel power gain from the \(i\)th beam over the \(k\)th carrier to the \(i\)th user and \(\sigma^2\) is the noise power. Hence, the offered capacity at beam \(i\) including all the possible available carriers is

\[
C_i = \sum_{k=1}^{K} B_{\text{sc}} \log_2(1 + \gamma_{i,k}(p_k, x_k)). \tag{2}
\]

Finally, given the \(i\)th beam demand \(D_i\), the normalized unmet system capacity is given by

\[
C_{\text{unmet}} = \sum_{i=1}^{N} \max(1 - C_i/D_i, 0). \tag{3}
\]

III. PROBLEM FORMULATION

In this section, we formulate the resource optimization problem to minimize the power and the bandwidth of the system while closely matching per user demand. In this context, we select the unmet system capacity, the total transmit power, and the total number of utilized carriers as multi-objective functions for our problem formulation. Furthermore, in order to account for the practical modulation and coding schemes employed in traditional satellite communication systems, we impose a minimum SINR constraint for all beams. In addition, we introduce a binary variable \(y_k\) to indicate whether \(k\)th carrier is active, i.e. assigned by any user \((y_k = 1)\) or not \((y_k = 0)\). The optimization problem is given by

\[
\begin{align*}
\text{minimize} & \quad \sum_{i=1}^{N} \max(1 - C_i/D_i, 0) + \sum_{k=1}^{K} y_k + \sum_{i=1}^{K} \sum_{k=1}^{K} y_k \sigma_k[i] \\
\text{subject to} & \quad T1: \gamma_{i,k}(x_k, \tilde{p}_k) \geq \gamma_{\text{min}}, \quad \forall i, \forall k, \\
& \quad T2: \sum_{k=1}^{K} \sum_{i=1}^{K} \tilde{p}_k[i] \leq 1, \\
& \quad T3: \sum_{k=1}^{K} \tilde{p}_k[i] P_{\text{total}} \leq P_{\text{max}}, \quad \forall i, \\
& \quad T4: y_k \in \{0, 1\}, \quad \forall k, \\
& \quad T5: x_k[i] \in \{0, 1\}, \quad \forall i, \forall k, \\
& \quad T6: x_k[i] \leq y_k, \quad \forall i, \forall k, \\
& \quad T7: P_{\text{total}} \tilde{p}_k[i] \leq x_k[i] P_{\text{max}}, \quad \forall i, \forall k, \\
& \quad T8: \tilde{p}_k[i] \geq 0, \quad \forall i, \forall k,
\end{align*}
\]

where \(\sum_{k=1}^{K} \gamma_k\) is the carrier utilization function and \(\sum_{i=1}^{N} \sum_{k=1}^{K} \tilde{p}_k[i]\) is the normalized power allocation function with \(\tilde{p}_k[i] = \frac{p_k[i]}{\sum_{k=1}^{K} p_k[i]}\). Note that we normalize the multi-objective function to avoid the mismatch unit and imbalance weight. The constraint \(T1\) imposes a minimum SINR \(\gamma_{\text{min}}\) needed per user. Constraint \(T2\) ensures that the sum-power consumption of all beams does not exceed the total power of the system \(P_{\text{total}}\). Moreover, the constraint \(T3\) imposes a limitation for the power allocation per beam to be within the maximal allowed beam power, i.e. \(P_{\text{max}}\). In addition, \(T4\) and \(T5\) indicate that the respective carrier activation variables \(x_k[i]\) and \(y_k\) are binary. \(T6\) indicates that the carrier assignment impacts the total number of active carriers. The constraint \(T7\) allow power to be allocated for \(x_k[i] = 1\). Finally, \(T8\) prohibits negative power allocation.

IV. DEMAND-AWARE POWER AND BANDWIDTH MINIMIZATION

The unmet system capacity, the non-linearity of \(\gamma_{i,k}(x_k, \tilde{p}_k)\) as well as the integer nature of \(x_k[i]\) and \(y_k\) makes the problem (4) non-convex. Therefore, problem (4) is a mixed-integer non-convex optimization problem, for which it is difficult to find the optimal solution. Furthermore, to find the optimal values of \(y_k\) and \(x_k\) an exhaustive search is required, which is computationally expensive even for problems with low dimension. Hence, to solve the problem, we first designed a suboptimal carrier allocation procedure where each carrier is allocated sequentially until the required capacity is provided to a user. Then, problem (4) with the above carrier allocation procedure reduces to a power allocation problem. We will show later how to address the power allocation problem. For \(x_k\) and \(y_k\), we apply the following carrier allocation procedure:

\[
\begin{align*}
y_m &= \begin{cases} 
0, & \sum_{k=1}^{m-1} C_{i,k} \geq D_i, \forall i \\
1, & \text{Otherwise}
\end{cases} \tag{5}
\end{align*}
\]

\[
\begin{align*}
x_m[i] &= \begin{cases} 
y_m, & \sum_{k=1}^{m-1} C_{i,k} < D_i, \\
0, & \text{Otherwise}
\end{cases} \tag{6}
\end{align*}
\]

\[
C_{i,k} = B_{\text{sc}} \log_2(1 + \gamma_{i,k}(\tilde{p}_k, x_k)), \tag{7}
\]

where (5) describes the \(m\)th carrier is available for transmission only if the previous \(m-1\) carriers are not enough to satisfy the demand. This happens when the offered capacity to any beam is below the required demand. Furthermore, (6) states that the \(m\)th carrier only available for beam \(i\) whenever the offered capacity is less than the requested demand. Then, using (5) and (6) as well as introducing an upper bound slack variable \(t_i\) to the unmet system capacity i.e. \(1 - C_i/D_i \leq t_i\), then the problem in (4) reduces to
The problem in (8) is a power allocation problem that can be solved iteratively using a quadratic transform, where the non-convex fractional part of the capacity is linearized [18]. Accordingly, the quadratic transform of $C_i$ is given by

$$
\tilde{C}_i(\gamma_{i,k}(\mathbf{x}_k, \mathbf{p}_k)) = \sum_{k=1}^{K} \log_2(1 + \gamma_{i,k}(\mathbf{x}_k, \mathbf{p}_k))
$$

(9)

$$
\gamma_{i,k}(\mathbf{x}_k, \mathbf{p}_k) = 2z_{i,k}\sqrt{g_{i,k}[j]x_k[i]p_k[j]P_{\text{total}}} - z_{i,k}^2 \left( \sum_{j=1,j\neq i}^{N} g_{i,k}[j]x_k[i]p_k[j]P_{\text{total}} + \sigma^2 \right)
$$

(10)

$$
z_k[i] = \frac{\sqrt{g_{i,k}[j]x_k[i]p_k[j]P_{\text{total}}}}{\sum_{j=1,j\neq i}^{N} g_{i,k}[j]x_k[i]p_k[j]P_{\text{total}} + \sigma^2}
$$

(11)

where $z_k[i]$ is the quadratic transform variable for each carrier and $\tilde{p}_k[i]$ is the previous value of $p_k[i]$. Finally, we re-write (8) as

$$
\begin{align*}
\text{minimize} & \quad \sum_{i=1}^{N} t_i + \sum_{i=1}^{N} \sum_{k=1}^{K} \tilde{p}_k[i] \\
\text{s.t.} & \quad T1, T2, T3, T7, T8, T10, \\
& \quad T9 : D_i - \tilde{C}_i(\gamma_{i,k}(\mathbf{x}_k, \mathbf{p}_k)) \leq t_i, \forall i,
\end{align*}
$$

(12)

The (12) is now convex and can be solved using the available tools of convex optimization, such as CVX [19]. The proposed solution is described using the algorithmic notation in Algorithm 1. First, we initially assign equal power to $\tilde{p}_k[i]$. Furthermore, we initialize $C_i = 0$, $y_k$ and $x_k$ to zero. Then, the algorithm starts to update $y_m$ and $x_m[i]$ according to (5) and then (6), respectively. Hence, for $m = 1$, the new updated value are $y_1 = 1$ and $x_1[i] = 1$. Then, the algorithm solves (12) and updates the value of $\tilde{p}_k[i]$ by the new values of $\tilde{p}_k[i]$. Subsequently, the algorithm updates the values of $z_k[i], y_m$ and $x_m[i]$ within the loop. The algorithm stops when the approximation of the quadratic form becomes accurate, i.e., the respective variable $z_k[i]$ does not change much within one iteration. Hence, the stopping condition is set as $\sum_{i=1}^{N} \sum_{k=1}^{K} |z_k[i] - z_k[i]| \leq 10^{-4}$, where $z_k[i]$ is the previous value of $z_k[i]$.

Algorithm 1: power and bandwidth allocation

Input: $C_i = 0$, $y_k = 0$, $x_k[i] = 0$, $m = 1$
update $y_m$ using (5) and $x_m[i]$ using (6);
repeat
$z_k[i] \leftarrow z_k[i]$; $\tilde{p}_k[i] = \frac{y_k[i]P_{\text{total}}}{\sum_{j=1}^{N} g_{i,k}[j]x_k[i]P_{\text{total}}}$; $z_k[i] = \sqrt{\frac{G_{i,k}[j]}{4\pi}}$
Solve (12); update $m$ using (5) and $x_m[i]$ using (6);
until $\sum_{i=1}^{N} \sum_{k=1}^{K} |z_k[i] - z_k[i]| \leq 10^{-4}$;
Output: $C_i$, $y_k$, $x_k[i], \forall i$, $\forall k$, $\tilde{p}_k[i] = \frac{y_k[i]P_{\text{total}}}{\sum_{j=1}^{N} g_{i,k}[j]x_k[i]P_{\text{total}}}$

V. Simulation Results

This section evaluates the performance of the proposed Demand-Aware Power and Bandwidth Minimization (DAPBM) scheme. The main simulation parameters are summarized in Table I. We consider $N = 100$ beams to cover a geographic area of Europe as shown in Fig. 1. The beam pattern is provided by the European Space Agency (ESA) and generated assuming a Direct Radiating Antenna (DRA), with 750 elements spaced by 5λ. Furthermore, the channel power gain for the $i$th user from the $k$th beam is given by $g_{i,k}[j] = \sqrt{\frac{G_{i,k}[j]}{4\pi}}$, where $G_{i,k}[j]$ is the user terminal antenna gain. $G_{i,k}[j]$ denotes the gain from the $i$th beam towards the $i$th user and $d_i$ is the slant range between the satellite and the $i$th user.

The results have been obtained with $M = 100$ Monte Carlo runs. In each run, a user location is selected randomly from the considered beam coverage area. An example of user locations for a single realization is shown in Fig. 1. For comparison, we considered reference [10] and [16] as benchmark schemes. In addition, we considered that all users have the same demand $D_i = D$, $\forall i$, and we vary the value of $D$ between 100 Mbps to 900 Mbps.

**Table I**

| Parameter | Value |
|-----------|-------|
| Satellite Orbit | 13°E |
| Number of beams ($N$) | 100 |
| Number of carriers ($K$) | 20 |
| System Bandwidth ($B_{\text{Total}}$) | 500 MHz |
| Sub-carrier bandwidth ($B_{\text{sub}}$) | 25 MHz |
| Noise power ($G_{\text{noise}}$) | 130 dBW |
| Minimum SINR ($\gamma_{\text{min}}$) | 2.2 dB |
| Max. beam gain ($G_{i,k}[j]$) | 51.3 dB |
| User antennas gain ($G_{i,k}$) | 9.8 dB |
| Total Power of the system ($P_{\text{Total}}$) | 1000W |
| Maximum power for each beam ($P_{\text{max}}$) | 100W |
A. Convergence analysis

The convergence of Algorithm 1 for different value of D is shown in Fig. 2. For the convergence analysis we observe the objective function of (12) i.e. \( F(t, \tilde{\beta}) = \sum_{i=1}^{N} t_i + \sum_{i=1}^{N} \sum_{k=1}^{K} \tilde{p}_k[i] \) at different iterations. We observed that the proposed Algorithm needs 6 iterations and 20 iterations to converge to the minimum value required to satisfy 100Mbps and 400Mbps demand per beam, respectively. Moreover, faster convergence is observed for lower demands in general, since fewer carriers are required.

B. Computational Complexity Analysis

The optimization (12) has \( K(N+1) \) decision variables and \( 2NK + 3N + 1 \) convex constraints. Assuming that the Algorithm 1 needs \( M \) iterations to converge. Hence, the approximate Big-O complexity is given by \( O(MK^3(N+1)(2NK + 3N + 1)) \approx O(MK^3) \) \cite{20}. Therefore, the Big-O of the proposed algorithm is a polynomial function that can efficiently run on computers in polynomial time. The actual computational time depends on the computer processors. For instance, using four 2.11 GHz cores and 16 GB of RAM, for a one-time resource allocation, the computational time for DAPBM, \cite{16}, and \cite{10} is 5 minutes, 4 minutes, and 90 minutes, respectively. Hence, the computational time of the proposed (12) is less than the exponential cone optimization used in \cite{10}. In contrast, \cite{16} has less computational time than DAPBM. However, the performance of \cite{16} is weaker than DAPBM as we will demonstrate in the next section.

C. Demand matching

In this section, we studied the demand satisfaction of the DAPBM and benchmark schemes. We compare the demand satisfaction using the following performance indicators

1) Average satisfaction index (ASI)

\[
ASI = \frac{1}{NM} \sum_{m=1}^{M} \sum_{n=1}^{N} \min \left( \frac{C_i}{D_m}, 1 \right)
\]

2) Average Unmet System Capacity (AUSC)

\[
AUSC = \frac{1}{NM} \sum_{m=1}^{M} \sum_{n=1}^{N} \max(D_i - C_i, 0)
\]

Fig. 3a provides the ASI for the DAPBM and the benchmark schemes. We observed decreased ASI from 450 Mbps for the proposed scheme and \cite{10}, whereas the starting decrease of ASI for \cite{16} is from 100 Mbps. Moreover, at 500 Mbps, the DAPBM provides 99% demand satisfaction but \cite{10} and \cite{16} have 97.5% and 89% demand satisfaction, respectively. Furthermore, on average, the DAPBM can satisfy the demand of 450 Mbps and 900 Mbps per beam at 100% and 60% cases, respectively but not in the benchmark schemes. Therefore, the proposed scheme outperforms the benchmark schemes. Moreover, in \cite{10}, the radiation power used in some carrier frequencies is small. Hence, poor SINR is observed below the system requirement. Consequently, lower demand satisfaction is observed. However, in DAPBM, a minimum SINR constraint is included to avoid the result of poor SINR. Hence, we gain more demand satisfaction compared to \cite{10}. For \cite{16}, an evolutionary algorithm is used to allocate the power and bandwidth, which can not guarantee the optimal solution. As a result, poor demand satisfaction is observed. In contrast, in DAPBM, higher demand satisfaction is obtained since it closely converges to the optimal solution.

Fig. 3b depicts the AUSC performance of the DAPBM and the benchmark schemes. The AUSC is zero below 450Mbps, for DAPBM and \cite{10}. In contrast, the AUSC is non zero for \cite{16}. Moreover, since the proposed algorithm minimizes the power and the carrier while closely matching the demand, it provides no excess offered capacity. Consequently, the mismatch error is zero up to 450 Mbps. However, we observed above 450 Mbps increase of the AUSC for all schemes. For instance, at 600 Mbps, the AUSC of DAPBM is 51 Mbps, whereas the AUSC of \cite{10} and \cite{16} is 61 Mbps and 132 Mbps, respectively. Hence, the system resource is not enough to accommodate higher demands. Nevertheless, precoding can be added on the top of the proposed scheme to improve demand satisfaction. Generally, the proposed scheme performed well for lower and medium demands.
D. Bandwidth and power consumption

In this section, we compare the bandwidth and the power consumption of all schemes. We use the following key performance indicators:

1) Average used power (AUP)

\[ \text{AUP} = \frac{1}{M} \sum_{m=1}^{M} \sum_{i=1}^{N} \sum_{k=1}^{K} p_k[i]. \]  

(15)

2) Average used bandwidth (AUB)

\[ \text{AUB} = \frac{B_{sc}}{M} \sum_{m=1}^{M} \sum_{k=1}^{K} y_k. \]  

(16)

Fig. 4a describes the AUB of all schemes. We have observed that the bandwidth utilization of the DAPBM is more efficient than the benchmark schemes. For instance, at 100 Mbps and 400 Mbps, respectively, the proposed scheme allocates 26% and 90% of the system bandwidth while [10] and [16] allocates 100% of the system bandwidth. Interestingly, the proposed algorithm allocates the bandwidth properly with respect to demand. For lower demands, a small portion of the bandwidth is required. As the demand increases, the bandwidth allocation increases. Therefore, the DAPBM allocates the system bandwidth dynamically with a change of per beam demand. This is not possible in the case of the benchmark schemes. However, for higher demand, the DAPBM usage of bandwidth converges to the system bandwidth.

The AUP in dBW of DAPBM and the benchmark schemes are shown in Fig.4b. We observed that the proposed scheme has less power consumption for low demand compared with the benchmark schemes. For example, at 100 Mbps the DAPBM is lower than [10] and [16] by 3 dBW and 2 dBW, respectively. For medium demands, the DAPBM used slightly more power compared to [10]. However, this more power consumption is compensated by saving more bandwidth, see Fig.4a. Though [16] have lower power consumption than the proposed scheme, the power allocation is not efficient. Consequently, poor demand satisfaction is observed. Generally, for higher demands, all scheme converges to the total power of the system.

E. Carrier allocation

In this section, we compare the carrier allocation of the DAPBM scheme at different demands. We use the following performance indicators:

1) Average used carriers (AUC)

\[ \text{AUC} = \frac{1}{M} \sum_{m=1}^{M} \sum_{k=1}^{K} y_k. \]  

(17)

2) Average unused carriers (AUNC)

\[ \text{AUNC} = \frac{1}{M} \sum_{m=1}^{M} \left( K - \sum_{k=1}^{K} y_k \right). \]  

(18)

Fig. 5 shows the AUC and AUNC of the proposed scheme. At low demand, the DAPBM utilizes fewer carriers. For instance,
at 100 Mbps, the AUNC is 15, and the AUC is 5. Furthermore, we observed that the AUC increases with respect to demand. However, the AUNC decreases when the demand increases. Generally, the proposed algorithm saves more unused carriers for demand below 400Mbps. Accordingly, those carriers can be applied to other services at the same time while satisfying the current system services. Finally, thanks to the proposed scheme, we can satisfy at most 400 Mbps per beam demand while saving 10% of the carriers.

Fig. 5. Carrier allocation of DAPBM

VI. CONCLUSIONS

In this paper, we proposed a novel method of Power and Bandwidth Minimization for Demand-Aware GEO Satellite Systems. The formulated optimization problem is non-convex in general. Hence, we use demand-aware carrier assignment with power allocation to solve the problem iteratively. The carrier assignment is done sequentially with the goal to not occupy additional carriers, if the demand is already satisfied. The power allocation remains a non-convex problem, which is solved using linearization and iterative update. The proposed scheme is compared with the benchmark schemes via extensive numerical simulations. We observed that the proposed scheme provides better performance in terms of demand satisfaction and bandwidth efficiency.

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