Abstract—Wireless sensor networks (WSNs) are used to monitor the environment where the networks are deployed. The lifetime of WSNs can be increased by energy-efficient or energy-balancing algorithms. Balanced energy consumption among all nodes is the main issue. In this paper, a new energy-efficient unequal clustering routing protocol (EEUCR) has been presented. In this protocol, the area of the network is divided into the number of rings of unequal size, and each ring is further divided into a number of clusters. Rings nearer to the base station (BS) have a smaller area and the area of rings keeps on increasing as the distance from BS increases. This helps to balance the energy consumption among the sensor nodes. The nodes with heterogeneous energy are deployed in the network. Nodes nearer to the base station have higher energy as compared to farther nodes. Static clustering is used but cluster heads (CHs) are not fixed and are elected on the basis of the remaining energy of the sensor node. Simulation results are compared with existing protocols and show improvement in energy consumption, which, in turn, increases the network lifetime of WSN and also balance the energy consumption of sensor nodes.

Keywords: energy balancing, hotspot problem, network lifetime, unequal clustering, wireless sensor networks.

I. INTRODUCTION

Advances in WSNs have provided tiny and less-costly sensors having the ability to sense various types of physical and environmental conditions[1][3]. They can process the data and can transmit that data to distant places through wireless communication. The main issue with WSN is that its energy depletes very hastily and becomes lifeless[18]. The maximum energy is consumed by the communication subsystem. To minimize energy consumption during communication, routing protocols can play a vital role. Therefore, to enhance the lifetime of the sensor node, routing algorithms should be energy-efficient[4]. The major goal of routing protocols is to find an energy-efficient and reliable route to transmit data from the sensor nodes to the sink so that the lifetime of the network is maximized[2]. The sensor nodes may communicate their data to BS through single-hop communication or multi-hop communication. In single-hop communication data is transmitted by every CH to the BS, this will lead to faster energy depletion of farther CHs as compared to CHs those are closer to the BS. This introduces the problem of the hotspot which results in network partitioning in an area that is farther from BS. But in multi-hop communication farther CHs send data to BS through hopping. This method is more energy-efficient than the single-hop communication but this introduces hotspot problems in the area nearer to the BS due to overburden of relay traffic.

Sensor nodes in the hotspot area exhaust their energy to a large extent which results in network partitioning near to the BS. So, there is a requirement of energy balancing among the nodes. A lot of work has been done in the last few years in the field of WSNs. We reviewed some of the most relevant papers. Heinzelman et al.[4] introduced clustering for the first time. This is a single-hop clustering algorithm. Adaptive clustering is used and also every node has the same probability to be elected as cluster head irrespective of remaining energy. TDMA (time division multiple access) is used for intra-cluster communication. Han et al.[5] proposed a hybrid unequal multi-hop clustering based on density (HCD) which enhances the lifetime of the network. This protocol selects the cluster head (CH) by comparing the status of each node to its adjacent nodes. It will consider a number of parameters for CH selection such as the energy of nodes, the number of adjacent nodes, the distance to the BS, and the layer where the node is placed. This technique tries to minimize the energy consumption of the network and improve load balancing. It follows unequal and hybrid dynamic–static clustering to achieve the objectives. Ye et al.[8] reported an energy-efficient clustering scheme. Cluster heads are elected with more residual energy through local radio communication while achieving well cluster head distribution. It also balances the load among the cluster heads. Yao et al.[9] focused on energy balancing to avoid network partitioning. The protocol is based on a heterogeneous network. The network is divided into rings of unequal sizes and every ring is further divided into clusters. As the rings approach BS, their size reduces. Two types of nodes are used regular sensor nodes (RSN) and CH. This avoids extra burden to elect the CH and CHs are of high energy than RSNs. Muhammad et al.[10] approached an energy-efficient hybrid routing protocol for underwater WSNs. In this protocol, WSN is divided into four regions of different radii. An equal number of nodes are deployed in all regions. Sensor nodes with more energy are deployed in region 1 and 2 as compared to sensor nodes deployed in the region 3 and 4 to balance the energy consumption. Smaragdakis et al.[12] proposed a protocol that uses the concept of the heterogeneous node to improve the stability period of the network. In this protocol, two types of nodes are deployed: advanced nodes and normal nodes. The advanced nodes have higher energy as compared to the normal nodes. Aderohunmu et al.[13] introduced the concept of three types of nodes in the network: advanced nodes, intermediate nodes, and normal nodes. This is also an improvement of SEP. The advanced nodes have higher energy than Intermediate and normal nodes. Intermediate nodes have energy level between advanced and normal nodes. Normal nodes have less energy as compared to other nodes. This method also improves life span of the WSN. Faisal et al.[14] describes an algorithm, in which sensor field is divided into three regions. Nodes in region near...
to the BS will directly send data to the BS. Sensor nodes deployed in the faraway region are higher energy nodes. These nodes follow clustering mechanism and data will be sent to BS by CH. Malathi et al.[15] highlighted a hybrid of static and dynamic clustering. In this algorithm network is divided into number of layers and then static clusters are formed. The parameters used to select the CH are residual energy of the sensor nodes, distance of CH from the sink and the number of their neighbors. Zhao et al.[16] illustrated a mechanism to manage the number of CH in the area where nodes are densely deployed. This algorithm improves the formula to find threshold by considering the distance between node and BS, the no. of nearest nodes, residual energy, and the average distance between sensor nodes. Khan et al.[17] proposed a cluster-based energy-efficient routing protocol. A CH receives data from root nodes, aggregate and then communicates aggregated data to the BS. This protocol uses static clustering. This protocol works well in both large and small areas. A large field is divided into rectangular clusters and then these clusters are grouped to form zones. This way communication between CH and BS is very efficient. Khan et al.[18] implemented a multistage data transmission mechanism. In this protocol CH is selected in a very efficient manner and static clustering is used which improves energy efficiency than dynamic clustering. Panag et al.[19] proposed dual-head static clustering algorithm to balance energy consumption by the sensor nodes. In this protocol two nodes in a cluster are selected as the CHs on the basis of remaining energy and distance from the BS and other nodes within the cluster. One CH is responsible for aggregation of data and second CH is responsible for transmission of data to BS. Gu et al. [20] have implemented a protocol that focuses on the coverage of the area of the WSN. The proposed protocol is ECDC that joins together both residual energy and the coverage importance metrics of the nodes to select the CH. Pramanick et al.[21] proposed a Clustering-based Routing scheme(EASSR) for WSNs. They use the sleep and wake property of some of the sensor nodes, which helps to increase the life of network. In this algorithm, a node is selected as a CH if its residual energy is greater than average energy of the system and if remaining energy of a node reaches five percent of its initial energy then sensor nodes can directly send data to BS.

After reviewing the literature it is found that a lot of work has been done for an energy-efficient network. Although many protocols proposed in literature try to use the energy of the network efficiently. Many researchers work to improve energy consumption among the nodes using unequal clustering[5-10] or deploying heterogeneous nodes[11-17] in the sensor field. Still, there is a requirement of some improved methods that can solve the problem in a better way.

In this paper, EEUCCR is proposed and it is a heterogeneous-aware unequal clustering routing protocol to solve the problem of the hotspot in WSNs and enhance the lifetime of the network. It is designed for a long-lived network and also enhances the time period before the exhaust of the first node of the network. It is assumed the network field is circular and BS with unlimited energy is placed in the middle of the field. The network is divided into five rings of unequal size around the BS and each ring is further divided into clusters as shown in fig 2. Also, the radii of rings and cluster size increases as their distance from the BS increases and vice-versa. This results in unequal clustering. The heterogeneous energy nodes are deployed in the network. The sensor nodes in rings 5, 4 and 3 are the same energy nodes. The energy of sensor nodes in rings 1 and 2 is two times extra than sensor nodes in rings 5, 4 and 3. The reason for increasing the energy of nodes in regions 1 and 2 is to balance the overburden of forwarding the relay traffic. This is beneficial for energy balancing among the network and tries to solve the hot spot problem and enhances network lifetime also shown by simulation results. Static clustering is used in this protocol. The remaining energy is considered as the parameter for the selection of CH. If the distance of CH to BS is less than d0 then there will be single-hop communication but if the distance is greater than equal to d0 then there will be multihop communication. The next-hop selection is done by considering two parameters: distance from current CH and remaining energy of the CHs in the lower ring.

The following sections discuss the remainder paper. Section II presents the network model. Section III shows the simulation results. In Section IV, conclusion and future scope is discussed.

### II. NETWORK MODEL

The EEUCCR uses the concept of heterogeneity of nodes, unequal clustering, and multi-hop communication. The protocol is divided into three different phases: initialization phase, setup phase, data transmission phase. Fig.1 shows the radio energy dissipation model[4]. The transmit electronics circuit will sense the data and then convert the data into the electronic form.

![Radio energy dissipation model](image)

**Fig.1. Radio energy dissipation model.**

The signal is amplified by the transmit amplifier and then transmitted by the antenna. Receiving antenna will receive the data and then processed by the receive electronics circuit. The following equations show the energy consumption to transmit and receive the k-bit packet to the distance (d) [4][12].

\[
E_{TX}(k,d) = E_{TX}(k) + E_{TX,amp}(k,d)
\]

\[
E_{TX}(k,d) = \begin{cases} 
E_{TX} * k + \epsilon_{fs} * k * d^2, & d < d_0 \\
E_{TX} * k + 2 \epsilon_{mp} * k * d^2, & d \geq d_0 
\end{cases}
\]

\[
E_{RX}(k) = E_{RX} * k
\]

Where, \[d_0 = \sqrt{\epsilon_{fs} / \epsilon_{mp}}\]

Equation 1 and 2, shows energy consumption for data transmission. where, \[E_{TX}(k)\] represent energy consumption for the transmission of k bits and \[E_{TX,amp}(k,d)\] is energy consumption for the amplification of k bits. Transmitter and receiver are apart by a distance(d). Equation 3, shows energy consumption for the reception of k bits. Energy consumption for the transmission increases with the increase of distance for the transmission. If \[d < d_0\] it will be single-hop communication but if \[d \geq d_0\] then this is multihop communication. The \[d_0\] as shown in equation 4, is critical distance and \[\epsilon_{fs} and \epsilon_{mp}\] are the constant coefficients for free-space communication and multipath communication,
respectively[4]. The value of $d_0$ depends upon system parameters such as the height of the antenna, the carrier signal wavelength, etc[9].

**A. Initialization Phase**

The initialization phase includes network architecture, node deployment, and clustering. This phase is necessary only one time in the origination before the communication rounds start. This phase includes the following steps.

**Step: 1 Network architecture**

In this model unequal clustering is proposed, we have considered that a BS is deployed at the center in the area under observation. The sensor field is divided into m rings. Radii of rings are not equal, this will help to balance the energy consumption of the sensor nodes. Radii reduce as the rings approach towards BS as shown in Fig 2 and Fig 3. radius $r_0<r_1<...<r_m$. The ring 0 has radius $r_0$, the ring 1 has radius $r_1$ and so on. Nodes those falls in radius (BS,$r_0$) will belong to ring 0, nodes those falls in radius ($r_0$, $r_1$) will belong to ring 1, it can also be stated that nodes in ($r_1$, $r_2$) will belong to ring i here, ($0 \leq i \leq m-1$). The radius of the innermost ring $r_0$ and the radius of the outermost ring $r_m$ is 0.5 [9], as it gives better results at this value. Radii of other rings are calculated by the following equation[9]:

$$r_m = r_0 + (r_m - r_0) \times i$$  \hspace{1cm} (5)

The following fig 2 shows Network architecture. This fig represents rings with different radii which results in unequal clustering.

![Network Architecture](image)

**Fig. 2. Network Architecture**

**Step: 2 Node deployment**

In this protocol, sensor nodes are deployed in a uniform random manner in WSN. BS is placed in the middle of the network field. Sensor nodes in the internal two rings $r_0$ & $r_1$ are higher energy nodes[10] and are called supernodes(Ns). The energy of nodes Ns is (E0) which is two times higher than the initial energy. Nodes in outer three rings $r_2$, $r_3$ & $r_4$ are of initial energy(E0) nodes and are called normal nodes(Nn). The total energy of the network is $T_E$ as given in equation (7). The area of the network is divided into two zones A0 and A0 where, A0 is an area where supernodes are deployed i.e. area of ring $r_0$ & $r_1$. A0 is an area where normal nodes are deployed i.e. area of rings $r_2$, $r_3$ & $r_4$. The number of nodes in zone A0 and A0 are proportional to their area and is given by the equation(6).

$$\frac{A_0}{A_n} = \frac{N_s}{N_n}$$  \hspace{1cm} (6)

$$T_E = E_0 + E_0$$  \hspace{1cm} (7)

**Step: 3 Clustering**

This protocol uses static clustering. Static clustering reduces the energy consumption than the dynamic clustering because it reduces extra communication required before every round in dynamic clustering. The algorithm divides the monitored area into predefined numbered clusters. Each node is aware of its position and transfers information to the BS through the control packet. The BS divides the network into four clusters per ring. Clusters are 90° apart from each other. The size of clusters is not the same in the network, it varies as the radii of rings varies.

**B. Setup Phase**

In the Setup phase, the proposed protocol selects the CHs and find the best route.

**Step: 1 CH selection**

The CH will be selected before the start of every communication round. The sensor node in the cluster having maximum residual energy will be selected as CH. A node in a cluster will compare its residual energy with other nodes in the cluster and after comparison node with higher residual energy will be selected as CH. This will maximize the life of the network because of the proper rotation of CH among a cluster.

**Step: 2 Route formation**

When a CH has been selected among the cluster, it will receive data from nodes of the cluster and other CHs and finally send the data to the BS. Let, CHj is cluster head of jth cluster and ith ring then it will calculate its distance from BS. If the distance of CHj to the BS is less than the critical distance then BS will be next-hop. Otherwise, it will use multi-hop communication and CH in the ring ($r_{i+1}$) will be selected as next-hop. This is shown in equation[19]

$$d_{CH_{i+1}} = \begin{cases} BS, & \text{if } d(CH_{i+1}, BS) < d_0 \\ CH_{i-1}, & \text{if } d(CH_{i+1}, BS) > d_0 \end{cases}$$  \hspace{1cm} (8)

When the $d$ is greater than $d_0$ then the next hop will be a CH in the lower-numbered ring and is selected by using the following steps.

When a CHj has to transfer data a list of negotiable-CHs(Nch) is prepared. An Nch is one that is closer to BS than the CHj or in the ring ($r_{i+1}$), and its distance must not more than $d_0$ from the CHj.

The fitness of the Nch is evaluated by the following equation:

$$F_{\text{Next-hop}}(N_{ch}) = w_1(E_{res,N_{ch}}) + w_2 \left( \frac{1}{d_{N_{ch},CH_j}} \right)$$  \hspace{1cm} (9)

Where $E_{res,N_{ch}}$ is the residual energy of Nch, $d_{N_{ch},CH_j}$ is the distance of Nch from the CHj. $w_1$ & $w_2$ are weighting factors and have a value between 0 & 1. The Nch has the highest value of $F_{\text{Next-hop}}$ is selected as the next-hop. After finding the best fitness, CHj will transfer the data to next-hop.

**C. Data transmission phase: Inter-cluster and intra-cluster communication**

Actual data communication takes place during this phase. The data transfer phase is divided into two parts–inter-cluster communication and inter-cluster communication. Intra-cluster communication means within the same cluster. In this phase, data is gathered by the cluster node in the $j$th cluster of the ring ($r_j$) and forwarded to the respective
CHs in their allocated time slot as per the TDMA schedule. TDMA is used here to avoid collision between the transmitted packets. During inter-cluster communication, CHs receive data from other CHs. The CHs aggregate the received data and then data is transmitted to the BS or to the next hop. The total energy consumed[12] in every cluster during each round is given in the following equation:

$$E_{\text{cluster}ij} = E_{\text{CH}ij} + lE_{\text{non CH}}$$  \hspace{1cm} (10)

The energy consumed by a sensor node which is not a CH, is given in the following formula:

$$E_{\text{non CH}} = k(E_{\text{TX}} + \epsilon f_s d_{ij}^2)$$  \hspace{1cm} (11)

$$d_{ij}^2$$ is the distance of a sensor node to its CH within the cluster.

If, \(l = \)the number of non- cluster head nodes in the cluster and \(i=m-1\), then \(E_{\text{CH}ij}\) is given as the following equation:

$$E_{\text{CH}ij} = E_{\text{RX}}(lk) + k(l+1)E_{\text{DA}} + k(l+1) (E_{\text{TX}} + \epsilon f_s d_{ij}^2)$$  \hspace{1cm} (12)

But, if \(0 \leq i \leq m-1\) then \(E_{\text{CH}ij}\) is given as the following equation.

$$E_{\text{CH}ij} = E_{\text{RX}}(\lambda_{ij} + lk) + (\lambda_{ij} + k(l+1))E_{\text{DA}} + (\lambda_{ij} + k(l+1)) (E_{\text{TX}} + \epsilon f_s d_{ij}^2)$$  \hspace{1cm} (13)

Here, \(E_{\text{DA}}\) is energy consumed for data aggregation. \(\lambda_{ij}\) is the bit arrival rate at \(CH_j\) from other \(CH\) of ring \(i-1\).

$$\lambda_{ij} = nk(l + 1), \quad n=\text{number of hops}$$  \hspace{1cm} (14)

### III. SIMULATION RESULTS AND DISCUSSION

The following parameters are measured by the simulations.

**Stability Period:** This is the time gap between the start of the operation of the network and the death of the first sensor node.

**Instability Period:** This is the time gap between the death of the first sensor node and the death of the last sensor node.

**Network lifetime:** This is the time between the start of the network operation and the death of the last node.

MatlabR2015b is used for simulations and is run over a computer with Intel® Core™ i5 CPU @2.4GHz having 4GB RAM. The following Table-I shows the value of parameters used for simulation. The protocol is executed many times to evaluate its performance. The details of the parameter values used in these scenarios are listed in Table-I. Three scenarios are taken to check the performance of EEUCR in different conditions. The EEUCR is compared with different groups of protocols in different scenarios because the results of the same protocols are not available in all the scenarios. Also, the total energy of EEUCR is equal to the total energy of the protocols used for comparison.

### Table –I: Parameters used for the simulation in three different scenarios.

| PARAMETERS                              | VALUES                  |
|-----------------------------------------|-------------------------|
| Network diameter                        | Scenario1                | Scenario2                | Scenario3                |
|                                         | 100m                    | 200m                    | 500m                    |
| Base station location                   | (50,50)                 | (100,100)               | (250,250)               |
| Number of nodes                         | 100                     | 100                     | 200                     |
| \(r_0\)                                | 50                      | 100                     | 250                     |
| \(r_1\)                                | 5                       | 10                      | 25                      |
| \(E_{\text{a}}\)                        | 0.5j                    | 0.82j                   | 0.41j                   |
| \(E_{\text{c}}\)                        | 1j                      | 1.64j                   | 0.82j                   |
| Total energy                            | 60j                     | 100j                    | 100j                    |
| Data packet size                        | 4000bits                | 8000bits                | 4000bits                |
| Number of rounds                        | 10000                   | 10000                   | 10000                   |
| \(m\)                                   | 5                       | 5                       | 5                       |
| \(q\)                                   | 0.5                     | 0.5                     | 0.5                     |
| \(E_{\text{DA}}\)                       | 5nj/bit                 | 5nj/bit                 | 5nj/bit                 |
| \(d_0\)                                 | 36.17m                  | 36.17m                  | 36.17m                  |
| Control packet size                     | 400bit                  | 400bit                  | 400bit                  |
| \(E_{\text{TX,REP}}\)                   | 0.014pj/bit             | 0.014pj/bit             | 0.014pj/bit             |
| \(E_{\text{TX}}\)                       | 50nj/bit                | 50nj/bit                | 50nj/bit                |
| \(E_{\text{RX}}\)                       | 50nj/bit                | 50nj/bit                | 50nj/bit                |
| \(E_{\text{a}}\)                        | 6.37x10^{-2}nj/bit/m^2  | 6.37x10^{-2}nj/bit/m^2  | 6.37x10^{-2}nj/bit/m^2  |
| \(E_{\text{a}}\)                        | 4.87x10^{-3}nj/bit/m^4  | 4.87x10^{-3}nj/bit/m^4  | 4.87x10^{-3}nj/bit/m^4  |

Following Fig.3 shows the formation of variable radii of rings as per scenario1. Ring formation for other scenarios will also be same. The only difference will be in their dimensions. The radii of rings increases as the distance from BS increases. The area is divided into five rings. The BS is located in the innermost ring of the network.
deployed in the area $A_n$ is approximately one ratio four as given in the following equation.

$$\frac{N_x}{N_{m}} = \frac{1}{4} \quad (15)$$

A. Energy consumption of EEUCR

As per scenario 1, the total energy of all the sensor nodes in the network is approximately 60j. The total energy of all the sensor nodes in the network as per scenario 2 is 100j and the total energy of network as per scenario 3 is also 100j. Fig. 4 shows energy consumption wrt round, graph 5(a), (b) and (c) shows energy consumption by the network as per scenario 1, 2 and 3, respectively. As shown in graph 5(a), high energy is consumed up to 1900 rounds because almost all the nodes are alive. After 3000 rounds, 80% of nodes die which results in less energy consumption after 3000 rounds. The graph 5(b) shows high energy is consumed up to 1600 rounds because a large number of nodes are alive. The graph 5(c) shows high energy is consumed up to 1000 rounds. A sheer slope of the graph shows that better energy balancing is achieved. The Fig. 5(a) for scenario 1 shows a sheer slope of the curve than the 5(b) and 5(c). So scenario 1 shows balanced energy consumption as compared to other scenarios.
Energy-Efficient unequal Clustering Routing Algorithm for WSNs

Following Fig. 6 shows the number of dead nodes w.r.t rounds and graphs 6(a), 6(b) and 6(c) shows dead node of the network as per scenario 1, 2 and 3, respectively.

As shown in graph 6(a), the stability period of the network is up to 1893 rounds. All the nodes die at 5024 rounds. The graph 6(b) shows the stability period of the network is up to 1061 rounds and all the nodes die at around 3541 rounds. The graph 6(c) shows the stability period of the network is up to 63 rounds and all the nodes die at 3780 rounds. Following Fig. 7 shows the number of alive nodes w.r.t rounds and graphs 7(a), 7(b) and 7(c) shows the alive node of the network as per scenario 1, 2 and 3, respectively. By comparing 7(a), 7(b) and 7(c), it is found that the instability period of the network is long in scenario 3 and also the life of the network is small than scenario 1. As shown in graph 7(a), the stability period and life of the network in scenario 1 is long as compared to scenario 2 and 3. The graph 7(b) shows that the stability period of the network is large than scenario 3 but smaller than scenario 1 and the life of the network is smaller than the other scenarios.
B. Comparison of the EEUCR with algorithms as per scenario 1

This section shows an evaluation of the EEUCR with other routing algorithms like LEACH[4], EASSR[15], SEP[12], ITSEP[16]. Table-II shows the values of first node dead (FND) and last node dead (LND) of all the protocols. In ITSEP, FND after 1400 rounds but 40% of nodes are alive after 2500 rounds but in the EEUCR only 30% of nodes are alive. The EEUCR shows better results of the stability period as compared to ITSEP. In LEACH the FND after 1022 and LND after 3500 rounds. In SEP, FND after 1158 rounds which is better than LEACH but the whole network dies after 2500 rounds. In EASSR, FND after 1350 rounds and LND after 4500 rounds. A sharp slope of EEUCR the graph indicates that better energy balancing is achieved until the first node dies. On reaching the FND round, the energy resources of many nodes start to deplete leading to the quick decline in the number of alive nodes. In EEUCR FND after 1893 rounds and network is alive up to 4544 rounds. Fig. 8 graphically shows a comparison of the number of nodes dead with respect to rounds for the existing routing algorithms with EEUCR.

C. Comparison of the EEUCR with algorithms as per scenario 2

This section shows an evaluation of the EEUCR with other routing algorithms HUCL[15], ECDC[20], HCD[5] for scenario 2. Results for comparison are taken from HCD[5] for all the protocols used for comparison in this scenario. Results of fig. 9 and table-III show that the routing algorithm EEUCR increases stability period and network lifetime than the other three protocols. In HUCL the FND after 200 rounds and LND after 1000 rounds. In ECDC, FND after 300 rounds which is better than HUCL and the whole network dies after 1400 rounds. In HCD, FND after 400 rounds and LND after 1200 rounds. In EEUCR, FND after 1061 rounds and network is alive up to 3780 rounds.

D. Comparison of the EBUCR with algorithms as per scenario 3

This section shows an evaluation of the EBUCR with other routing algorithms for large scenario 3. AZR-LEACH[17] is static clustering routing algorithms and results for all the protocols used for comparison in this scenario are taken from this paper. The LEACH, SEP are dynamic clustering routing algorithms. Results show that the stability period of the EEUCR is slightly smaller than AZR-LEACH but greater than LEACH & SEP. The network lifetime and instability period of the EBUCR are greater than all the protocols as shown in Fig. 10 & Table-IV. In LEACH the FND after 39 rounds and LND after 1078 rounds. In SEP, FND after 62 rounds and the whole network dies...
after 1282 rounds. In AZR-LEACH, FND after 80 rounds and LND after 1400 rounds. In EEUCR, FND after 63 rounds and whole network dies after 3541 rounds. In this scenario EEUCR performs better than all protocols used for comparison in terms of lifetime of WSN. Table-IV shows that value of FND of EEUCR over the FND of LEACH is 1.61, which is second highest in the table. It is also shown that value of LND of EEUCR over the LND of LEACH is 3.28, which is larger than all the algorithms used for the comparison.

IV. CONCLUSION AND FUTURE SCOPE

The proposed protocol is an unequal clustering based routing protocol to balance the energy consumption among nodes. The network area is divided into rings of unequal size so that nodes nearer to BS have less diameter and can save energy. Heterogeneous nodes are deployed in the sensor field. In the ring 1 and 2 supernodes are deployed and in all other rings, normal nodes are deployed. This hybrid network reduces the overhead to relay traffic of sensor nodes of rings 1 and 2. This helps to balance the energy consumption among the network which leads to energy efficiency and further increases the network lifetime. After simulation results of EEUCR are compared with existing protocols in three different scenarios. Almost in all scenarios proposed protocol shows better performance than the existing protocols. The authors would be evaluating the performance of the EEUCR for the packet delivery ratio and end to end propagation delay in their future work.


table I:

Table-II: FND and LND values of LEACH, EASSR, SEP, ITSEP, and EEUCR for scenario1.

| Protocol  | FND    | LND    |
|-----------|--------|--------|
|           | Round  | Round  |
| LEACH     | 1022   | 3500   |
| EASSR     | 1350   | 4500   |
| SEP       | 1158   | 2500   |
| ITSEP     | 1400   | 40%    |
| EEUCR     | 1893   | 1200   |

Table –III: FND and LND values of HUCL, ECDC, HCD, and EEUCR for scenario2.

| Protocol  | FND    | LND    |
|-----------|--------|--------|
|           | Round  | Round  |
| HUCL      | 200    | 1000   |
| ECDC      | 300    | 1400   |
| HCD       | 400    | 1200   |
| EEUCR     | 1061   | 3780   |

Table –IV: FND and LND values of LEACH, AZR-LEACH, SEP, and EEUCR for scenario3.

| Protocol  | FND    | LND    |
|-----------|--------|--------|
|           | Round  | Round  |
| LEACH     | 39     | 1078   |
| SEP       | 62     | 1282   |
| AZR-LEACH | 80     | 1400   |
| EEUCR     | 63     | 3541   |

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AUTHORS PROFILE

Priyanka Honda received her B.E. in Electronics and Communication Engineering from MDU, Rohtak in 2006 and M.E. in Electronics and Communication Engineering from Panjab University, Chandigarh in 2010. She is pursuing her Ph.D. from Punjab technical university, Jalandhar. At present she is working as an Assistant Professor in the department of Electronics and Communication Engineering at Ch. Devi Lal State institute of Engg. And Technology,panniwala mota, Haryana, India. She is having 10 years of teaching and research experience. Her areas of research is Wireless Sensor Networks and Optimization techniques. She is a life member of Indian Society for Technical Education (ISTE).

Tripatjot Singh Panag obtained his Bachelor’s degree in Electronics and Communication Engineering from Guru Nanak Dev Engineering College (GDSEC), Ludhiana, Punjab, in May 2001 and Master’s degree in Electronics and Communication Engineering from Baba Banda Singh Bahadur Engineering College (BBSBEC), Fatehgarh Sahib, Punjab, in April 2008. He completed his Ph.D. from Sant Longowal Institute of Engineering and Technology, Longowal, Punjab in November 2017. At present he is working as an Assistant Professor in the department of Electronics and Communication Engineering at Baba Banda Singh Bahadur Engineering College (BBSBEC), Fatehgarh Sahib, Punjab, India. He is having 16 years of teaching and research experience. His areas of interest are Wireless Communication, Data Communication Networks, Digital System Design, Wireless Sensor Networks and Optimization techniques. He is a life member of Indian Society for Technical Education (ISTE).

Balwinder Singh Sohi received his B.Sc. Engineering, Masters of Engineering and Ph.D. in Electronics from Panjab University, Chandigarh, India in 1971, 1981 and 1992 respectively. He is the former Director of University Institute of Engineering and Technology, Panjab University. He is recipient of ‘The Sir Thomas Ward Memorial Prize’ of Institute of Engineers, Kolkata, in 2004 and 2005. He is a senior member of IEEE, fellow of Institution of Engineers, Life Member of ISTE. Presently he is working as director at Chandigarh Group of Colleges, Gharuan, Punjab. He has more than hundred research publications in various reputed journals and conferences. His areas of research are microprocessors and their applications, biomedical instrumentation, wireless networking, computer networking, microwave and RF engineering.