Comparation and Analysis of Performance Parameters of Basic Adders with Sparse Adder

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Abstract—Adders play a vital role in the design of a digital system using VLSI (Very Large Scale Integration) technique. Adders are the basic building block of ALU (Arithmetic Logic Unit) which is an important component of a processor. In this paper we are comparing and analyzing the performance parameters of basic adders like Ripple Carry Adder, Carry Select Adder, Carry Look Ahead Adder, Parallel Prefix Adder along with sparse adder. The above mentioned adders are implemented using 90nm technology in Xilinx ISE 14.7 Suite.

Keywords—Ripple Carry Adder, Carry Select Adder, Carry Look Ahead Adder, Parallel Prefix Adder, Xilinx ISE 14.7 Suite.

I. INTRODUCTION

Addition is the most important operation of any digital circuit. Adders not only perform addition but also performs some basic arithmetic operations like subtraction, multiplication, increment and decrement. Adders form the heart of the digital circuit that improves the performance of the ALU which in turn increases the efficiency of the processor.

II. TYPES OF ADDERS

A. HALF ADDER

Half adder is a combinational arithmetic circuit that adds two numbers and produces a sum bit (S) and carry bit (C) as the output. If A and B are the input bits, then sum bit (S) is the X-OR of A and B and the carry bit (C) will be the AND of A and B. Thus Half adder circuit is constructed using one X-OR gate and one AND gate. The half adder can add only two input bits (A and B) and has nothing to do with the carry if there is any in the input.

B. FULL ADDER

Full Adder is the adder which adds three inputs and produces two outputs. The first two inputs are A and B and the third input is an input carry as C-IN. The output carry is designated as C-OUT and the normal output is designated as S which is SUM.
C. Ripple Carry Adder
The Ripple carry adder is constructed by cascading full adder blocks in series. Each full adder is responsible for adding two binary digits at any stage. The carry-out of one stage is directly fed as the carry-in to the next stage. Each stage has to wait for the carry generated from the previous stage for computation.

| A | B | Cin | Sum | Cout |
|---|---|-----|-----|------|
| 0 | 0 | 0   | 0   | 0    |
| 0 | 0 | 1   | 1   | 0    |
| 0 | 1 | 0   | 0   | 1    |
| 0 | 1 | 1   | 0   | 1    |
| 1 | 0 | 0   | 1   | 0    |
| 1 | 0 | 1   | 0   | 1    |
| 1 | 1 | 0   | 0   | 1    |
| 1 | 1 | 1   | 1   | 1    |

Truth table of Full adder

D. Carry Select Adder
The Carry select adder pre-computes the sum and carry-out for two possible cases of input carry that is when Cin = 0 and Cin = 1. The calculated sum is given to the multiplexer, which chooses the correct output depending upon the input carry coming from previous stage. This pre-computation of Sum reduces the delay of rippling of carry which is limited to only one multiplexer for each stage.

E. Carry Look Ahead Adder
 Carry look-ahead adder is designed to overcome the latency introduced by the rippling effect of the carry bits. This adder is based on the principle of looking at the lower order bits of the augends and addend if a higher order carry is generated. In this adder the outputs Sum and Carry are derived from the intermediate terms defined as Generate (G) and Propagate (P) terms.

Generate term: \( G = A \cdot B \)
Propagate term: \( P = A \oplus B \).

F. Parallel Prefix Adder
The Parallel prefix adder consists of three stages: pre-computation stage, prefix network stage and post-computation stage. The pre-computation stage computes the carry ‘Propagate’ and carry ‘Generate’ bits for each input pair. The prefix network stage computes the final carries from the carry ‘Propagate’ and carry ‘Generate’ bits. The final post computation stage computes the final Sum from carry generated in prefix network stage.
Sparse adders are highly efficient adders and are majorly known for their high speed computations. In this paper we have constructed the sparse adder as a combination of Kogge-Stone adder which is a form of parallel prefix adder along with Carry Look Ahead adder.

III. COMPARISON ANALYSIS

Comparing the basic adders like Ripple Carry Adder, Carry Select Adder and Carry Look Ahead Adder with the performance parameters: delay, power, number of bonded IOBs, number of occupied slices, total number of 4 input LUTs.

Table 1: Comparison of basic adders with performance parameters

| ADDER                  | DELAY (ns) | POWER (mW) | NO: OF BONDED IOBs | NO: OF OCCUPIED SLICES | TOTAL NO: OF 4 INPUT LUTs |
|------------------------|------------|------------|--------------------|------------------------|--------------------------|
| RIPPLE CARRY ADDER     | 12.008     | 68         | 14                 | 6                      | 8                        |
| CARRY SELECT ADDER     | 10.739     | 76         | 14                 | 6                      | 9                        |
| CARRY LOOK AHEAD ADDER | 11.625     | 62         | 14                 | 7                      | 11                       |

Comparing the parallel prefix adder and sparse adder with the performance parameters: delay, power, number of bonded IOBs, number of occupied slices, total number of 4 input LUTs.

Table 2: Comparing parallel prefix adder and sparse adder with performance parameters

| ADDER                        | DELAY (ns) | POWER (mW) | NO: OF BONDED IOBs | NO: OF OCCUPIED SLICES | TOTAL NO: OF 4 INPUT LUTs |
|------------------------------|------------|------------|--------------------|------------------------|--------------------------|
| PARALLEL PREFIX ADDER        | 26.012     | 193        | 49                 | 22                     | 42                       |
| SPARSE ADDER                 | 16.816     | 164        | 49                 | 39                     | 73                       |

IV. RESULT

In this paper sparse adder is implemented and analysed. According to the observations made sparse adder is efficient in terms of delay and power, also wiring complexity of the sparse adder is greatly reduced when compared with parallel prefix adder.
V. SIMULATIONS

Simulation of 4-bit Ripple Carry Adder

Simulation of 4-bit Carry Select Adder

Simulation of 4-bit Carry Look Ahead Adder
VI. CONCLUSION
In this paper we have implemented Ripple Carry Adder, Carry Select Adder, Carry Look Ahead Adder, Parallel Prefix Adder and Sparse Adder using 90nm technology in Xilinx ISE 14.7 Suite. From the comparison analysis we can conclude that sparse adder is efficient in terms of both power and delay. The delay is less for sparse adder when compared to parallel prefix adder. Power is more for carry select adder and less for sparse adder. Since sparse adder occupies more area, it can be optimized or reduced by exploring certain algorithms.

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