Imperfect and Unmatched CSIT is Still Useful for the Frequency Correlated MISO Broadcast Channel

Chenxi Hao and Bruno Clerckx
Communication and Signal Processing Group, Department of Electrical and Electronic Engineering
Imperial College London, United Kingdom
Email: {chenxi.hao10,b.clerckx}@imperial.ac.uk

Abstract—Since Maddah-Ali and Tse showed that the completely stale transmitter-side channel state information (CSIT) still benefits the Degrees of Freedom (DoF) of the Multiple-Input-Multiple-Output (MISO) Broadcast Channel (BC), there has been much interest in the academic literature to investigate the impact of imperfect CSIT on DoF region of time correlated broadcast channel. Even though the research focus has been on time correlated channels so far, a similar but different problem concerns the frequency correlated channels. Indeed, the imperfect CSIT also impacts the DoF region of frequency correlated channels, as exemplified by current multi-carrier wireless systems.

This contribution, for the first time in the literature, investigates a general frequency correlated setting where a two-antenna transmitter has imperfect knowledge of CSI of two single-antenna users on two adjacent subbands. A new scheme is derived as an integration of Zero-Forcing Beamforming (ZFBF) and the scheme proposed by Maddah-Ali and Tse. The achievable DoF region resulted by this scheme is expressed as a function of the qualities of CSIT.  

I. INTRODUCTION

In downlink multi-user multiple-input-multiple-output (MU-MIMO) communications, the latency and inaccurate CSIT degrade the DoF when conventional precoding techniques such as ZFBF are employed. Strategies to exploit imperfect feedback to enhance DoF region has therefore attracted a lot of attention. The completely stale CSIT was first studied by Maddah-Ali and Tse. In their contribution [1], an optimal per-user DoF of \( \frac{1}{2} \) in a two-user setup was achieved by a simple transmission scheme (denoted as MAT scheme in the sequel). That result was extended later in [2] and [3], by accounting for imperfect current CSIT. The optimal DoF was derived and expressed as a function of the quality of the current CSIT. The achievability was shown using a scheme that bridges the DoF found in [1] and [2][3].

In [4] and [5], unequal quality of current CSIT per user was investigated. The optimal bound found is a superset of the results in [2] and [3], revealing that an asymmetric DoF is achieved by each user. Moreover, [6] has studied the imperfect delayed CSIT, suggesting that it can be used as good as perfect delayed CSIT.

To date, all the works only focus on the time correlated channel. However, the DoF region of frequency correlated channels is also impacted by the imperfect CSIT. In current multi-carrier communication systems, the CSIT is measured and reported by users on a per-subband basis. In practice, each user only reports its CSI on a group of predefined subbands, which might provide few information about the channel of the subbands outside the group because of the weak correlation between different subbands. In this paper, we investigate, for the first time in the literature, the DoF of a general two-subband based frequency correlated broadcast channel with arbitrary imperfect CSIT (see Section II). Our contributions are summarized as follows:

1) Derive an achievable DoF of a two-user and two-subband based scenario as a function of the quality of the CSIT;
2) Design a novel transmission strategy, motivated by MAT and ZFBF, that achieves the DoF region.

The rest of this paper is organized as follows. The system model is introduced in Section III and the achievable DoF region is given in Section IV. The DoF achieved via using MAT and ZFBF is identified in Section V and a novel transmission scheme is introduced. Section VI concludes the paper.

The following notations are used throughout the paper. Bold lower letters stand for vectors whereas a symbol not in bold font represents a scalar. \((\cdot)^T\) and \((\cdot)^H\) represent the transpose and conjugate transpose of a matrix or vector respectively. \(\mathbf{h}^\perp\) denotes the orthogonal space of the channel vector \(\mathbf{h}\). \(\mathcal{E}[\cdot]\) refers to the expectation of a random variable, vector or matrix. \(|\cdot|\) refers to the magnitude of a scalar. \(f(P) \sim P^\alpha\) corresponds to \(\lim_{P \to \infty} \frac{\log(f(P))}{\log(P)} = \alpha\), where \(P\) is supposed to be the SNR throughout the paper and logarithms are in base 2. \(P_\alpha\) denotes the power of \(a\) while \(R^1_a\) and \(R^2_a\) represent the rate of \(a\) achieved at receiver 1 and 2 respectively.

II. SYSTEM MODEL

We consider a two-user broadcast channel with two transmit antennas and one antenna per user. The related parameters are defined as follows. \(h_i\) and \(g_i\) are the channel states in subband \(i\) of user 1 and user 2 respectively. Denoting the transmit signal vector in subband \(i\) as \(s_i\), subject to a per-subband based power constraint \(\mathcal{E}[|s_i|^2] \sim P\), the observations at receiver 1 and 2,
The error vectors are in the first subband using pilots and feeds it back as $h^\varepsilon$. As shown, we assume the qualities of $h_1$ with perfect knowledge of local CSI; 5) It is important to note the quantities $\mathbb{E}\left[||h_i^H g_i||^2\right]=\mathbb{E}\left[||g_i^H h_i||^2\right]\sim P^{-\beta}$ and $\mathbb{E}\left[||g_i^H g_i||^2\right]=\mathbb{E}\left[||h_i^H h_i||^2\right]\sim P^{-\alpha}$.

$y_i$ and $z_i$, respectively, can be written as

$$y_i = h_i^H s_i + \epsilon_{i,y}, \quad (1)$$

$$z_i = g_i^H s_i + \epsilon_{i,z}, \quad (2)$$

where $\epsilon_{i,y}$ and $\epsilon_{i,z}$ are unit power AWGN noise. Signal vector $s_i$ is a function of the symbol vectors for user 1 and user 2, denoted as $u_i$ and $v_i$, respectively. $u_i$ is a two-element symbol vector containing $u_{i,1}$ and $u_{i,2}$. Similarly, $v_i$ is defined.

The channels are characterized as follows. $h_i$ and $g_i$ are mutually independent and identically distributed with zero mean and unit covariance matrix ($\mathbb{E}\left[||h_i^H g_i||^2\right]=0$ and $\mathbb{E}\left[||h_i h_i^H||\right]=I_2$). The imperfect CSIT of user 1 is denoted as $\hat{h}_i$ while the imperfect CSIT of user 2 is $\hat{g}_i$, each with the error vector of $h_i = h_i - \hat{h}_i$ and $g_i = g_i - \hat{g}_i$. The variances of the error vectors are $\mathbb{E}\left[||\hat{h}_i||^2\right]=\sigma^2_{h,i}$ and $\mathbb{E}\left[||\hat{g}_i||^2\right]=\sigma^2_{g,i}$.

The CSIT setting in this two-subband based scenario is illustrated in Figure 1. User 1 estimates its channel information in the first subband using pilots and feeds it back as $\hat{h}_1$ while user 2 reports its CSI in the second subband as $\hat{g}_2$. As shown, we assume the qualities of $h_1$ and $g_2$ are identical and expressed using a parameter, $\beta$, which is defined as

$$\beta \equiv \lim_{P \rightarrow \infty} \frac{-\log \sigma^2_{h,1}}{\log P} = \lim_{P \rightarrow \infty} \frac{-\log \sigma^2_{g,2}}{\log P}. \quad (3)$$

As the CSI of two adjacent subbands are correlated, the transmitter can predict the channel information of the unreported subband. To be specific, with the knowledge of $\hat{h}_1$, the channel condition of the second subband of user 1, $\hat{h}_2$, is predicted. Similarly, $\hat{g}_1$ is predicted based on the knowledge of $\hat{g}_2$. The qualities of these two predicted channel states are characterized as $\alpha$, which is defined as

$$\alpha \equiv \lim_{P \rightarrow \infty} \frac{-\log \sigma^2_{h,2}}{\log P} = \lim_{P \rightarrow \infty} \frac{-\log \sigma^2_{g,1}}{\log P}. \quad (4)$$

Remarks: 1) $\beta$ and $\alpha$ vary within the range of $[0,1]$, where 0 represents no CSIT whereas 1 stands for perfect CSIT; 2) The quality of the predicted CSI, $\hat{h}_2$ and $\hat{g}_1$, is bounded by the quality of $h_1$ and $g_2$, namely $\alpha \leq \beta$; 3) We assume that this two-subband scenario can be repeated an infinite number of times; 4) The transmitter and both users have the knowledge of $\hat{h}_{1:2N}$ and $\hat{g}_{1:2N}$. Besides, each receiver has perfect knowledge of local CSI; 5) It is important to note the quantities $\mathbb{E}\left[||h_i^H h_i||^2\right]=\mathbb{E}\left[||g_i^H g_i||^2\right]\sim P^{-\beta}$ and $\mathbb{E}\left[||g_i^H g_i||^2\right]=\mathbb{E}\left[||h_i^H h_i||^2\right]\sim P^{-\alpha}$.

Throughout the paper, we define a per-channel-use based DoF, which is expressed as

$$d_i \triangleq \lim_{P \rightarrow \infty} \frac{R_i}{S \log P}, \quad i = 1,2, \quad (5)$$

where $R_i$ is the rate achieved by user $i$ over $S$ channel uses.

### III. DoF Region with General CSIT Pattern

**Theorem 1.** In a frequency correlated MISO BC with imperfect CSIT shown in Figure 1, an achievable DoF region is characterized as a polygon composed of the following corner points:

$$\{ (1,0) ; (0,1) ; (1,\alpha) ; (\alpha,1) ; \ldots \}$$

$$\{ \min \left( \frac{2+\alpha}{3}, \beta \right), \max \left( \frac{2+\alpha}{3}, \frac{2-\beta+\alpha}{2} \right) ; \ldots \}$$

$$\{ \max \left( \frac{2+\alpha}{3}, \frac{2-\beta+\alpha}{2} \right), \min \left( \frac{2+\alpha}{3}, \beta \right) \} \} \quad (6)$$

Figure 2 illustrates the region specified by (6), spanning all $\beta$ and $\alpha$ satisfying $\beta, \alpha \in [0,1]$ and $\alpha \leq \beta$. When $\alpha$ is fixed, points C and D (shown by circle points, see Figure 2) move closer to each other as $\beta$ increases. For $\beta=\frac{2+\alpha}{3}$, points C and D join at point E (or F, see Figure 2). If $\beta$ continues increasing, the DoF region will not expand any further. This reveals that the CSIT with quality $\beta$ satisfying $\beta \geq \frac{2+\alpha}{3}$ can be as good as $\beta=1$. Specifically, the DoF region achieved by MAT with $\beta=1$ and $\alpha=0$ (composed of point F, (1,0) and (0,1), see Figure 2) can be actually achieved by $\beta=\frac{2+\alpha}{3}$ and $\alpha=0$.

Moreover, if we fix $\beta$ and increase $\alpha$, all the points will move either upwards or to the right. When $\alpha$ reaches $\beta$, point A will join point C while points B and D overlap. The DoF region can be simply achieved by doing ZFBF plus superposition coding.

$^2$S channel uses may physically refer to $S$ subbands with full transmission power $P$ or transmitting symbols using power of $P^S$ in a single subband. In this paper, it takes the latter understanding.
In addition, the maximum sum DoF is achieved by the diamond point E, and star point F, for the case $\beta \geq \frac{2 + \alpha}{3}$. Otherwise, it is obtained by the circle points C and D.

IV. ACHIEVABILITY

A. Motivations

In this part, we briefly revisit two existing schemes, MAT and ZFBF. Their achievable rates in frequency correlated BC will be identified and analyzed. Their sub-optimalitys will motivate the derivation of the novel transmission strategy.

1) Reusing MAT scheme and Extensions: In [1], the transmission of MAT finishes in three time slots, during which the transmit signal and received signals are

$$s_1 = v_1, \quad s_2 = u_2, \quad s_3 = [\eta_{1,1} + \eta_{2,2}, 0]^T,$$

$$y_1 = \eta_{1,1}, \quad y_2 = h_{1}^H u_2, \quad y_3 = h_{2,1}^H (\eta_{1,1} + \eta_{2,2}),$$

$$z_1 = g_{1}^H v_1, \quad z_2 = \eta_{2,2}, \quad z_3 = g_{2,1}^H (\eta_{1,1} + \eta_{2,2}),$$

where $\eta_{1,2} = h_{1}^H v_2$ and $\eta_{2,2} = g_{2}^H u_2$. User 1 receives its desired symbol vector $u_2$ in $y_2$ but overhears $h_{1}^H v_1$ in $y_1$. The decoding is enabled once the transmission at slot 3 is completed, where the sum of the overhead interference is retransmitted. After decoding $\eta_{1,1} + \eta_{2,2}$ received in $y_2$ and subtracting $y_1$, user 1 obtains an additional independent observation of its desired symbol vector, $g_{2}^H u_2$. Hence, user 1 can decode $u_2$. Similarly, user 2 can decode $v_1$. In this way, four symbols are successfully transmitted in three slots, resulting in the symmetric DoF of $\frac{4}{3}$.

However, among all the six CSI in these three time slots, only two of them are in fact employed, namely $h_{1}^H$ and $g_{2}^H$. Equivalently, we can reuse MAT in the scenario shown in Figure 1 provided that $\beta = 1$ and $\alpha = 0$. The sum of the overhead interference, $\eta_{1,1} + \eta_{2,2}$, is reconstructed and retransmitted using an extra channel use (subband 3). The CSI of this extra channel use does not have to be known at the transmitter.

When $\beta < 1$, the transmit power should be adjusted because the overhead interferences generated at each user are restructured with non-negligible error at the transmitter. Specifically, after subtracting $y_1$ from $y_3 = h_{3,1}^H g_{2}^H u_2$, we obtain a residue interference, $\tilde{h}_{1}^H - h_{1}^H u_2$, obtained plus a residue interference, $\tilde{h}_{1}^H - h_{1}^H v_1$, where

$$\mathcal{E} \left[ \| \tilde{h}_{1}^H \|_2 \right] \sim P^{-\beta}. \quad (8)$$

To make the residue interference drown by the noise, the transmission power of $v_1$ in subband 1 should be reduced to $P^\alpha$. In this way, $\beta$ channel use is employed per subband, during which both $v_1$ and $u_2$ achieve the rate $2\beta \log P$ resulting in the sum DoF $\frac{4\beta}{3}$ over $3\beta$ channel uses.

2) Conventional interference mitigation techniques that achieve MUMIMO transmission. The transmitter precodes two symbols $u_1$ and $v_1$ (intended to user 1 and 2 respectively) using the knowledge of CSIT of both users. The transmission signal in subband 1 is expressed as

$$s_1 = \tilde{g}_{1}^H u_1 + \tilde{h}_{1}^H v_1, \quad (7)$$

where $P_{u_1} \sim P^\alpha$ and $P_{v_1} \sim P^\beta$, resulting in the received signals

$$y_1 = h_{1}^H \tilde{g}_{1}^H u_1 + h_{1}^H \tilde{h}_{1}^H v_1 + \epsilon_1, \quad (8)$$

$$z_1 = g_{1}^H \tilde{g}_{1}^H u_1 + g_{1}^H \tilde{h}_{1}^H v_1 + \epsilon_2. \quad (9)$$

As the qualities of $\tilde{h}_{1}^H$ and $\tilde{g}_{1}^H$ are $\beta$ and $\alpha$ respectively, $h_{1}^H \tilde{h}_{1}^H v_1$ and $g_{1}^H \tilde{g}_{1}^H u_1$ are drowned by the noise. In this way, the rate achieved by $u_1$ and $v_1$ are $\alpha \log P$ and $\beta \log P$ respectively. The amount of channel use in subband 1 is $\beta$. Similarly, the same transmission is applicable in subband 2 by switching the power of each user’s symbol. Hence, the sum DoF is $\frac{2\beta + 2\alpha}{3}$ during these $2\beta$ channel uses.

| Sum Rate (log $P$) | Channel Uses | $n_{1,1}, n_{1,2}$ | $n_{2,1}, n_{2,2}$ | Preceding |
|--------------------|--------------|--------------------|--------------------|-----------|
| MAT $4\beta$       | $3\beta$     | $2\beta$           | $1\beta$           | No        |
| ZFBF $2\beta + 2\alpha$ | $2\beta$ | $1\beta$           | $1\beta$           | Yes       |
| Objective $4\beta$ | $< 3\beta$   | $1\beta$           | $2\beta$           | Yes       |

TABLE I: Comparison among MAT, ZFBF and the objective. $n_{i,j}$ refers to the number of private symbols sent to user $i$ in subband $j$.

3) Analysis and motivation: The comparison between MAT and ZFBF are presented in Table II. ZFBF saves $\beta$ channel uses while it incurs a rate loss of $2(\beta - \alpha) \log P$ compared to MAT.

When $\alpha$ is small, MAT outperforms ZFBF in sum DoF. In this case, ZFBF precoding works inefficiently in rejecting the interference potentially seen by user 2 in subband 1, the transmit power of $u_1$ in (7) is therefore significantly limited, resulting in low DoF. Similarly, user 2 achieves low rate in subband 2. However, MAT transmits two symbols to each user in turn. The CSIT with quality $\beta$ is exploited to provide confident side information over an extra channel use. The DoF is therefore boosted up.

When $\alpha$ approaches $\beta$, ZFBF works well in rejecting the interference potentially seen by both users in each subband. The sum rate achieved by ZFBF therefore approximates as $4\beta \log P$, resulting in a higher sum DoF than MAT by saving the $\beta$ extra channel uses. However, MAT incurs a loss because the CSIT with quality $\alpha$ is wasted during the $3\beta$ channel uses.

Intuitively, given a certain value of $\alpha \in [0, \beta]$, a better sum DoF can be obtained by a strategy that optimally balances the employment of CSIT with quality $\alpha$ and the usage of extra channel use. This objective strategy can be designed as the integration of ZFBF and MAT. It would outperform ZFBF by employing a small fraction of extra channel use to perform overhead interference cancellation. At the same time, as precoding is introduced, the amount of extra channel use could be reduced compared to MAT while the sum rate remains $4\beta \log P$. The amount of extra channel use would be a function of $\beta$ and $\alpha$, bridging ZFBF and MAT. When $\alpha = \beta$, the transmission scheme would be upgraded to ZFBF; for $\alpha = 0$, it would collapse to pure MAT. Bearing this in mind, we derive the novel transmission block in the following section.

B. Building New Transmission Blocks

Following the aforementioned motivation, the main features of this strategy are presented in the last row of Table II. It is
a combination of ZFBF and MAT in terms of precoding and the number of symbols transmitted to each user per subband. The transmission signals in subband 1 and 2 are respectively expressed as

\[ s_1 = [x_{c,1}, 0]^T + [\mu_1, 0]^T + \left[ \mathbf{h}_1^T, \mathbf{g}_1^T \right]^T v_1 + \mathbf{g}_1^T u_1, \]

\[ s_2 = [x_{c,2}, 0]^T + [\mu_2, 0]^T + \left[ \mathbf{h}_2^T, \mathbf{g}_2^T \right]^T u_2 + \mathbf{h}_2^T v_2, \]

where two private symbols \((v_1 = [v_{1,1}, v_{1,2}]^T)\) are transmitted to user 2 and one private symbol \((u_1)\) is sent to user 1 in subband 1. Precoding is also considered in \(s_1\), where \(v_{1,2}\) is precoded with \(\mathbf{h}_1\), \(v_{1,1}\) and \(u_1\) are projected to the orthogonal space of \(\mathbf{h}_1\) and \(\mathbf{g}_1\) respectively. The new symbols in \(s_2\) are similarly encoded and transmitted.

Besides, \(\mu_1\) and \(\mu_2\) are two pieces of the quantized overheard interference, encoded with the rates, \(R_{\mu_1} = r_{\mu_1} \log P\) and \(R_{\mu_2} = r_{\mu_2} \log P\) respectively. \(x_{c,1}\) and \(x_{c,2}\) are common messages that should be decoded by both users. The power and rate allocation are presented in Table II.

| Subband 1 | Symbols | Power | Encoding rate \((\log P)\) |
|-----------|---------|-------|--------------------------|
| \(x_{c,1}\) | \(P - P^{\mu_1 + \beta}\) | \(1 - r_{\mu_1} - \beta\) |
| \(\mu_1\) | \(P^{\mu_1 + \beta} - P^\beta\) | \(r_{\mu_1}\) |
| \(v_{1,1}\) | \(P^\beta/2 - P^\alpha/2\) | \(\beta - \alpha\) |
| \(v_{1,2}\) | \(P^\beta/2 - P^\alpha/2\) | \(\beta - \alpha\) |
| \(u_1\) | \(P^\beta/2 - P^\alpha/2\) | \(\beta - \alpha\) |

| Subband 2 | Symbols | Power | Encoding rate \((\log P)\) |
|-----------|---------|-------|--------------------------|
| \(x_{c,2}\) | \(P - P^{\mu_2 + \beta}\) | \(1 - r_{\mu_2} - \beta\) |
| \(\mu_2\) | \(P^{\mu_2 + \beta} - P^\beta\) | \(r_{\mu_2}\) |
| \(v_{2,1}\) | \(P^\beta/2 - P^\alpha/2\) | \(\beta - \alpha\) |
| \(v_{2,2}\) | \(P^\beta/2 - P^\alpha/2\) | \(\beta - \alpha\) |

TABLE II: The power and rate allocated to the symbols in (10) and (11).

The received signals at each receiver in each subband are expressed as

\[ y_1 = h_{1,1}^* x_{c,1} + h_{1,1}^* \mu_1 + \eta_{1,1} + h_{1,1}^* \eta_{1,2} + v_1 + \epsilon_{1,1}, \]

\[ z_1 = h_{1,1}^* x_{c,1} + g_{1,1}^* \mu_1 + h_{1,1}^* \eta_{1,2} + v_1 + g_{1,1}^* u_1 + \epsilon_{1,2}, \]

\[ y_2 = h_{2,1}^* x_{c,2} + h_{2,1}^* \mu_2 + h_{2,1}^* \eta_{2,2} + u_2 + h_{2,1}^* v_2 + \epsilon_{2,1}, \]

\[ z_2 = g_{2,1}^* x_{c,2} + g_{2,1}^* \mu_2 + \eta_{2,2} + g_{2,1}^* v_2 + \epsilon_{2,2}, \]

where \(\eta_{1,1}\) and \(\eta_{2,2}\) are the overheard interferences generated at user 1 in subband 1 and at user 2 in subband 2 respectively,

\[ \eta_{1,1} = h_{1,1}^* \eta_{1,1} v_{1,1} + h_{1,1}^* v_{1,2}, \]

\[ \eta_{2,2} = g_{2,2}^* u_{2,1} + g_{2,2}^* u_{2,2}. \]

Note that the power stated below each term is obtained asymptotically, which is merely valid at high SNR. Since \(\mathbb{E}\left[|h_{1,1}^* h_{1,1}^*|^2\right] \sim P^{-\beta}\), the term \(h_{1,1}^* h_{1,1}^* v_{1,1}\) in \(\eta_{1,1}\) is drowned by the noise in (12). Similarly, \(g_{2,2}^* g_{2,2}^* u_{2,1}\) in \(\eta_{2,2}\) vanishes.

Hence, the overheard interference \(\eta_{1,1}\) and \(\eta_{2,2}\) are composed of \(v_{1,2}\) and \(u_{2,2}\) respectively, which are then possible to be detected at receiver 1 and 2 respectively (will be discussed in Section V). In this way, when reconstructing the sum of the overheard interferences, the channel component, for instance, \(h_{1,1}^* h_{1,1}^*\) in (16), can be dropped. As a consequence, in contrast to MAT where \(\eta_{1,1}\) is reconstructed and sent, we reconstruct the sum of the symbols \(v_{1,2}\) and \(u_{2,2}\) as

\[ \mu = v_{1,2} + u_{2,2}. \]

\(\mu\) can be generated from a codebook \(\{v_{1,2}, u_{2,2}\}\). Since \(v_{1,2}\) and \(u_{2,2}\) are encoded with same rate, we assume they are generated from the same codebook, denoted as \(\Phi\). Moreover, we design \(\Phi\) as a set close to the arithmetic plus \(\mathbb{Z}\). In this way, \(\mu=v_{1,2}+u_{2,2}\) can be generated from \(\Phi\) as well, with the encoding rate, \(r_{\mu} = \beta - \alpha\), identical to that for \(v_{1,2}\) and \(u_{2,2}\).

Furthermore, as motivated by (4) and (6), we split \(\mu\) into two parts, \(\mu_1\) and \(\mu_2\) as

\[ \mu = \{\mu_1, \mu_2\}, \]

each is encoded from a codebook \(\Phi_1\) and \(\Phi_2\) respectively, which are independent to each other. The encoding rates of these two codebooks are subject to

\[ R_{\mu_1} + R_{\mu_2} = (\beta - \alpha) \log P, \]

\[ r_{\mu_1} + r_{\mu_2} = \beta - \alpha. \]

Hence, \(\Phi\) can be considered as a product set of \(\Phi_1\) and \(\Phi_2\), \(\mu_1\) and \(\mu_2\) are decoded separately in two parallel channels, \(\mu\) can therefore be perfectly reconstructed by combining them.

As presented in Table III \(\mu_1\) and \(\mu_2\) are respectively superposed on the private symbols transmitted in subband 1 and 2. However, their power \(P^{\mu_1 + \beta}\) and \(P^{\mu_2 + \beta}\) should not exceed the power constraint \(P\). This constraint can be expressed as

\[ r_{\mu_1} \leq 1 - \beta, \quad r_{\mu_2} \leq 1 - \beta. \]

As a consequence, the transmission is subject to the relationship between (21) and (22).

First, in the case of \(2(1-\beta) > \beta - \alpha\), namely \(\beta < \frac{2}{\alpha+\beta}\), the power of \(\mu_1\) and \(\mu_2\) does not exceed the per-subband power constraint by simply setting \(r_{\mu_1} = r_{\mu_2} = \frac{2}{\alpha+\beta}\). Moreover, we can superimpose a common message \(x_{c,1}\) on \(\mu_1\) in subband 1 and \(x_{c,2}\) on \(\mu_2\) in subband 2 using power stated in Table III which is scaled with \(P\). Second, when \(2(1-\beta) = \beta - \alpha\), the power constraint is still satisfied but no common messages is transmitted since \(r_{\mu_1} + r_{\mu_2} = 1\). Third, for the case of \(2(1-\beta) < \beta - \alpha\), namely \(\beta > \frac{2}{\alpha+\beta}\), the value of \(r_{\mu_1}\) and \(r_{\mu_2}\) are bounded by \(1-\beta\) as in (22). Therefore, \(\mu\) has to be divided into three pieces as

\[ \tilde{\mu} = \{\mu_1, \mu_2, \mu_3\}. \]

\(\Phi\) close to arithmetic plus is defined as, if \(a \in \Phi\) and \(b \in \Phi\), then \(a + b \in \Phi\).
with the rates given by
\[ R_{\mu_1} = R_{\mu_2} \approx \log P^{r_{\mu_1}}, \quad R_{\mu_3} \approx \log P^{r_{\mu_3}}, \] (24)
\[ r_{\mu_1} = r_{\mu_2} = 1 - \beta, \quad r_{\mu_3} = 3\beta - \alpha - 2. \] (25)
The transmission of \( \mu_3 \) requires an extra channel use in another subband. Next, we will discuss the achievable rates of each point in Figure 2 depending on the requirement of extra channel use.

### C. Case I: \( \beta \leq \frac{2\alpha}{3} \) Achieving Points C and D

In this case, \( \mu \) is split into two parts and no extra channel use is required. Messages \( x_{c,1} \) and \( x_{c,2} \) are transmitted provided that \( \beta < \frac{2\alpha}{3} \). The decoding procedure is described as follows.

1) **Stage I**: Decode \( x_{c,1} \) and \( x_{c,2} \): Revisiting (12) and (13), the received power of \( x_{c,1} \) is \( P \) at each receiver. Successive interference cancelation (SIC) is selected as the decoding strategy. \( x_{c,1} \) is decoded at the first stage treating all the other symbols as noise. Consequently, the rates of \( x_{c,1} \) achieved at user 1 and user 2 are
\[ R_{x_{c,1}} = I(x_{c,1}; y_1 | h_1) \] and \( R_{x_{c,2}} = I(x_{c,1}; z_1 | g_1) \), respectively. These two rates are equal to \( \log P^{1-r_{\mu_1}-\beta} \) which is asymptotically \( (1-r_{\mu_1}-\beta) \log P \) for infinite \( P \). Similarly, \( x_{c,2} \) achieves the rate \( (1-r_{\mu_2}-\beta) \log P \) in \( y_2 \) and \( z_2 \).

2) **Stage II**: Decode \( \mu_1, \mu_2 \) and obtain \( \hat{\mu} \): As \( \mu_1 \) and \( \mu_2 \) are independently encoded and sent in subband 1 and 2 respectively, they can be decoded separately at both receivers. After that, \( \mu \) is obtained by combining them.

In \( y_1 \) and \( z_1 \), \( \mu_1 \) is decoded at the second stage of SIC, where \( x_{c,1} \) has been completely transmitted. Treating all the component to the r.h.s. of \( \mu_1 \) in (12) and (13) as noise, \( \mu_1 \) is decoded with the rate \( R_{\mu_1} = I(\mu_1; y_1 | h_1, x_{c,1}) \) and \( R_{\mu_1} = I(\mu_1; z_1 | g_1, x_{c,1}) \) by user 1 and user 2 respectively. Both quantities are equal to \( \log P^{1-r_{\mu_1}} \), which is \( r_{\mu_1} \log P \) at high SNR. Similarly, \( \mu_2 \) is decoded with rate \( r_{\mu_2} \log P \) in subband 2. After that, \( \mu_1 \) and \( \mu_2 \) have been successfully decoded so that \( \mu \) is completely recovered.

3) **Stage III**: Decode \( u_1 \) and \( u_2 \): Employing SIC as the decoding strategy, \( u_1 \) and \( u_2 \) are decoded from \( y_1 \) and \( z_2 \) respectively.

Let us introduce a notation, \( y_{1}' \), representing the signal after subtracting \( x_{c,1} \) and \( \mu_1 \) as
\[ y_1' = y_1 - h_{1,1}(x_{c,1} + \mu_1) \] (26)
\[ = \eta_{1,1} + h_{1,1}^H \sigma_{1,1} u_1 + \epsilon_{1,y} \] (27)
\[ = h_{1,1}^H \sigma_{1,1} u_1 + h_{1,1}^H \sigma_{1,1} u_1 + \epsilon_{1,y} \] (28)
where \( \eta_{1,1} \) is given in (16) and \( \epsilon_{1,y} \) results from merging \( \epsilon_{1,y} \) and \( h_{1,1}^H \sigma_{1,1} u_1 \). Treating \( h_{1,1}^H \sigma_{1,1} u_1 \) as noise, \( v_{1,2} \) can be decoded with the rate \( R_{v_{1,2}} = I(v_{1,2}; y_1 | h_1, h_2) \) = \( \log P^{r_{\mu_1}+p_{\mu_1}} \), which is asymptotically equal to \( (1-\beta-\alpha) \log P \) at high SNR. After that, \( u_1 \) is seen by subtracting \( h_{1,1}^H \sigma_{1,1} u_1 \) from \( y_1' \) as
\[ y_1'' = y_1' - h_{1,1}^H \sigma_{1,1} u_1 = h_{1,1}^H \sigma_{1,1} u_1 + \epsilon_{1,y}. \] (29)
The rate of \( u_1 \) is \( R_{u_1} = I(u_1; y_1'' | h_1, g_1) = \alpha \log P \). Similarly, employing SIC to \( z_2 \) results in \( R_{v_2} = \alpha \log P \) and \( R_{u_{2,2}} = (\beta-\alpha) \log P \).

### D. Case II: \( \beta > \frac{2\alpha}{3} \) Achieving Point E

In this case, we remind the reader of the discussion in Section IV-B that \( \mu \) is split into three parts and an extra channel use is required to transmit \( \mu_3 \). Besides, no common message is transmitted.

To achieve point E, we repeat the transmission blocks in (10) and (11) for \( L \) times and employ one additional subband, namely subband \( 2L+1 \), to finalize the transmissions of \( \mu_{3,i}, i=1,2,\ldots,L \), where \( \mu_{3,i} \) refers to the third piece of overhead interference generated in subband \( 2i-1 \) and \( 2i \). The rate of \( \mu_{3,i} \) is denoted as \( r_{\mu_{3,i}} \log P \) and we assume \( r_{\mu_{3,i}} = r_{\mu_{3,i}} = 1,2,\ldots,L \). The quality of CSIT in subband \( 2L+1 \) is identical to subband 1.

The transmission in subband \( 2L+1 \) is expressed as
\[ s_{2L+1} = [\mu_{3,1,0}]^T + [\mu_{3,2,0}]^T + \ldots + [\mu_{3,L,0}]^T + \ldots \] (33)
\[ g_{2L+1}^H u_{2L+1} + h_{2L+1}^H v_{2L+1}, \]
with the power and rate allocation presented in Table I
Considering \( s_{2L+1} \) and the transmit power, the received signal at user 1 is given as
\[ y_{2L+1} = h_{2L+1,1}^H \mu_{3,1} + h_{2L+1,2}^H \mu_{3,2} + \ldots + h_{2L+1,L}^H \mu_{3,L} + \ldots \]
\[ 1 - \beta \] (34)
where all the symbols are decodable using SIC. Specifically, after \( \mu_{3,1,i-1,1} \) are decoded, \( \mu_{3,i} \) are decoded treating all the
components to the r.h.s. of it in (33) as noise. The rate achieved for $\mu_{3,i}, i < L$ is derived as

$$R_{\mu_{3,i}}^{(1)} = I (\mu_{3,i}; y_{2L+1} | h_{2L+1}, \mu_{3,1:i-1})$$

$$= \log P_{\mu_{3,i}} - \log \left( \sum_{j=1}^{i-1} P_{\mu_{3,j}} \right) \geq \log \left( \frac{p^{1-(i-1)r_{\mu_{3,i}}} - p^{1-r_{\mu_{3,i}}}}{1 - \alpha} \right)$$

After decoding $\mu_{3,1:L-1}$, $\mu_{3,L}$ can be decoded treating $u_{2L+1}$ as noise. The rate of $\mu_{3,L}$ is $R_{\mu_{3,L}}^{(1)} = \log \left( \frac{p^{1-(L-1)r_{\mu_{3,L}}} - p^{1-r_{\mu_{3,L}}}}{1 - \alpha} \right)$, whose pre-log factor is $1 - (L-1)r_{\mu_{3,L}} - \alpha$. To make $\mu_{3,L}$ decodable with rate $r_{\mu_{3,L}}$, $L$ should satisfy the condition $1 - (L-1)r_{\mu_{3,L}} = r_{\mu_{3,L}} - \alpha$, resulting in

$$L = \frac{1}{\alpha} - \frac{1}{\beta} = \frac{1 - \alpha}{3\beta - 2}.$$  

Similarly, user 2 can decode $\mu_{3,1:L}$ using SIC.

Consequently, $\mu_{1:2:L}$ can be recovered by collecting and combining $\mu_{1,1:L}$, $\mu_{2,1:L}$ and $\mu_{3,1:L}$. Moreover, all the symbols transmitted from the 1st to 2Lth subband are decodable using the decoding flow described in Section IV-C. The rates achieved in subband $2i-1$ and $2i$ are

$$R_{u_{2i-1}}^{(2)} = \log P_{\mu_{3,L}}, R_{u_{2i}}^{(2)} = \log P_{\mu_{3,L}}.$$  

Besides, $u_{2L+1}$ is decoded in (33) with rate $R_{u_{2L+1}}^{(2)} = \log P_{\mu_{3,L}}$ at the last stage of SIC after removing all the $\mu_{3,i}$. Similarly, $R_{u_{2L+1}}^{(2)} = \alpha$. $R_{u_{2L+1}}^{(2)}$ is achieved at receiver 2. Finally, we can conclude the DoF achieved by each user as

$$d_1 = \lim_{P \to \infty} \frac{L \left( R_{u_{2i-1}}^{(1)} + R_{u_{2i}}^{(1)} + R_{u_{2L+1}}^{(1)} \right)}{(2L+1) \log P} = \frac{2 + \alpha}{3}, \quad (39)$$

$$d_2 = \lim_{P \to \infty} \frac{L \left( R_{u_{2i-1}}^{(2)} + R_{u_{2i}}^{(2)} + R_{u_{2L+1}}^{(2)} \right)}{(2L+1) \log P} = \frac{2 + \alpha}{3}.$$  

In the proposed scheme, the transmissions of private symbols in subband 1 and 2 occupy $2\beta$ channel uses while transmitting $\hat{\mu}$ requires $\beta - \alpha$ channel uses. Over those channel uses, the sum rate of the private symbols are $2\beta \log P$, resulting in the sum DoF $\beta \alpha$. Revisiting Table I our new scheme achieves the same sum rate as MAT but using less channel uses. At the same time, it outperforms ZFBF by $2(\beta - \alpha) \log P$ in sum rate with only $\beta - \alpha$ more channel use.

E. "SC+ZF"-Achieving Points A and B

Points A and B can be simply achieved via ZFBF using $\alpha$ channel use and transmitting common message $x_{c,i}$ using the remaining $1 - \alpha$ channel use in each subband. The transmitted signal in subband 1 is expressed as

$$s_1 = [x_{c,1}, 0]^T + g_{0,1}^T u_1 + h_{0,1}^T v_1,$$

where $u_1$, $v_1$ are allocated with power $P^\alpha/2$ and the power of $x_{c,1}$ is $P - P^\alpha$. User 1 observes $x_{c,1}$ plus $h_{0,1}^H g_{0,1}^T u_1$ because $h_{0,1}^H h_{0,1}^T v_1$ is drowned by the noise. Treating $h_{0,1}^H g_{0,1}^T u_1$ as noise, $x_{c,1}$ is decoded at the first stage of SIC with rate $(1 - \alpha) \log P$. After that, $u_1$ is decoded with the rate $\alpha$. Similarly, user 2 can decode $x_{c,2}$ and $x_{c,2}$ are intended to user 1, the DoF are expressed as

$$d_1 = \lim_{P \to \infty} \frac{R_{x_{c,1}}^{(1)} + R_{x_{c,2}}^{(1)} + R_{u_1}^{(1)} + R_{u_2}^{(1)}}{2 \log P} = 1,$$

$$d_2 = \lim_{P \to \infty} \frac{R_{x_{c,2}}^{(2)} + R_{x_{c,2}}^{(2)}}{2 \log P} = \alpha,$$

so that point B is achieved. Point A is achieved if the common messages are assumed intended to user 2.

V. CONCLUSION

This work for the first time in the literature investigates the impact of imperfect CSIT on the DoF region of frequency correlated MISO BC. A general two-subband based imperfect CSIT pattern (see Figure I) is studied. MAT and ZFBF achieve the optimal sum DoF for $\alpha = 0$ and $\alpha = \beta$ respectively while both of them have DoF loss for $0 < \alpha < \beta$. A novel transmission strategy is proposed to improve the performance. It processes as an integration of MAT and ZFBF, where precoding and overhead interference cancellation are combined. The DoF region achieved is a function of $\beta$ and $\alpha$, enhancing the result for $0 < \alpha < \beta$ and bridging the region of MAT (for $\alpha = 0$) and ZFBF (for $\alpha = \beta$).

More details on the achievability and the converse will be provided in our upcoming journal version paper. Besides, we will investigate a more general scenario where the qualities of the four CSIT in Figure I are all different. The proposed scheme and the DoF region will be extended.

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