Efficient Maintenance Scheduler for Near Optimum Utilization of Oil Tanks

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Abstract: Due to the entry of Kuwait into WTO and the expected liberalization of petroleum management, the oil market is to become more competitive. However, the space limitation and public increasing awareness of environmental protection issues and stricter regulation passed by Environmental Public Authority (EPA) in Kuwait make the need for full utilization of oil tanks a prime requirement of a successful oil business. In order to help oil companies to achieve this and maximize revenues by increasing the availability of tanks, an efficient maintenance scheduling is needed. This study introduces a new hybrid evolutionary algorithm and its implantation for solving real-world problem of oil storage tanks. The algorithm incorporates the American Petroleum Institute (API) standard 650 for open inspection procedure to produce a near optimal schedule for maintenance though continuously preserving a population diversity that ensures solution quality and convergence efficiency. The computational results show that the proposed evolutionary algorithm outperforms existing scheduler in literature and produces a higher quality solution which means better revenues though more operational tanks and more environmental protection though maximizing the utilization of existing tanks and less new tanks to build.

Key words: Evolutionary, maintenance scheduling, environmental protection, increase revenue, Oil, Kuwait

INTRODUCTION

In accordance with the entry of Kuwait into WTO entry and the recently issued Petroleum Management Law, the barrier of oil market in Kuwait will been removed. International petroleum vendors can sell their oil products in Kuwait current dominating markets in not so far future. This liberalization results in high competition in the oil market. To keep a cutting-edge position, competitors have to efficiently and effectively control the distribution channel of products, which consists of gas stations, pipelines and storage tanks. Due to the space limitation and the residents’ increasing awareness in environmental protection issues and stricter regulation passed by Environmental Public Authority (EPA) in Kuwait, the construction of storage tanks is the toughest obstacle for Kuwaiti and international vendors. In addition, according to the Petroleum Management Law, refining vendors or importers must always maintain reserves of 60 days or 50,000 kiloliters. Therefore, they unavoidably have to optimize the tank uses from the domestic oil companies. See Fig. 1 for an example of oil tank farms in Kuwait.

On the other hand, following the American Petroleum Institute (API) standard 650, storage tanks must be inspected every two years. Depending upon the corrosion degree inside the tanks, a so-called “open inspection” procedure will be conducted every five -to ten years. Each tank will take 60 to 240 days of outage for open inspection based on different capacity and construction type. As a result, a well-devised maintenance schedule of storage tanks will substantially help the leaseholder increase revenue attributed to the availability of tanks but assure the statutory reserves without constructing new tanks.

Currently, there is one domestic oil vendor in Kuwait, Kuwait Oil Company (KOC). In the past, the maintenance scheduling relies on the tacit knowledge of senior engineers or the package of linear programming. For a larger number of tanks, the increasing complexity is too high to be managed manually. In the literature,
genetic algorithms (GA) have been shown able to tackle complicated scheduling problems\cite{1-7}. In particular, GA outperforms other heuristic search approaches, such as simulated annealing and tabu search due to the fact (that it is relatively easy to encode in heuristic space and problem space\cite{8}). However, GA is subject to suffer the problem of premature convergence, which makes it tend to fall into local optimum. To amend such limitation, we propose to apply the hybrid evolutionary algorithm (heuristic)\cite{9} to deal with this real-world scheduling problem.

**Problem statement:** The process of petroleum refinement automatically runs 24 hours a day. At first the refined oil is stored in storage tanks and then transported from tanks to gas stations through pipelines or by tank trucks. According to API STD 650, the storage tanks must go through a periodic open inspection. The objective is to find a satisfying maintenance scheduling for the outage caused by open inspection in one year. In the refining system, the outage of storage tanks will affect the stability of oil supply. The level of effect is determined by the net reserve of the oil tank company. The net reserve in certain month \( m \) is defined:

\[
N_m = C - \sum_{i=1}^{T} \eta_{i,m} - \omega_m
\]

\( N_m \) the obtained net reserve in month \( m \),
\( C \): the total capacity of oil Tank Company,
\( \eta_{i,m} \): the capacity of outage of the \( i \)-th tank in month \( m \),
\( T \): the number of tanks and
\( \omega_m \): the forecasting maximum load in month \( m \).

In addition, there are two constraints for this maintenance scheduling problem:
* The process of maintenance must begin on the first day of a month and end on the last day of a month. Furthermore, the maintenance should be on schedule and cannot be abandoned.
* The volume of net reserve must be greater than zero at any time. The objective is to keep the net reserve maximum during maintenance.

The period of maintenance scheduling under investigation is one year in this study; that is to say, there will be 12 monthly net reserves in one year. Based on conservative estimation, we determined the lowest net reserves as the fitness of the schedule. i.e.

\[
\text{Net} = \min N_m = \min \left\{ C - \sum_{i=1}^{T} \eta_{i,m} - \omega_m, m = 1 - 12 \right\}
\]

**THE PROPOSED HEURISTICS**

**Simulate Annealing Heuristic (SA):** The main idea behind simulated annealing heuristic is to have a large number of iterations where in every iteration of the heuristic there is a single random pair exchange in the sequence. If the exchange improves the objective function then it accepts the exchange and the new sequence is preserved. If the objective function does not improve, then it is only allowed to accept the exchange with some small probability \( p \). As the number of iterations increases, the probability \( p \) for which the heuristic is allowed to accept an exchange that does not improve the objective function is reduced exponentially. This reduction in the probability usually is expressed as a function of a start temperature that is reduced by a cooling factor to reach a final (freezing) temperature. This technique of reducing the probability of accepting non-improving exchanges has proven to be very useful in escaping local optima’s during the course of search for global optima. The following is an algorithmic description of the heuristic.

**Simulate Annealing Heuristic (SA)**

\[
\text{Begin} \quad \text{Let } T_1 = 0.1 \quad \text{While } T_1 \geq 0.0001 \quad \text{Begin} \quad \text{Repeat 50 times} \quad \text{Begin} \quad \text{Let } L_1 = \text{value of the objective function} \quad \text{with current sequence} \quad \text{Pick two random positions } j \text{ and } k \quad \text{Swap jobs in the positions of } j \text{ and } k \quad \text{Let } L_2 = \text{value of the objective function} \quad \text{after the swap} \quad \text{If } L_2 < L_1 \text{ then accept swap} \quad \text{If } L_2 > L_1 \text{ then accept the move with} \quad \text{probability } f \text{ where} \quad d = \frac{L_2 - L_1}{L_1} \quad f = e^{-\frac{d}{T_1}} \quad \text{End Repeat} \quad \text{Let } T_1 = T_1 * 0.98 \quad \text{End While}
\]

**End Heuristic**

Setting the parameters for the proposed simulated annealing heuristic is essential in achieving a good performance. After some experimentation, the parameters for the simulated annealing heuristic are set as follows; the initial temperature \( T_1 \) is set to 0.1, the cooling factor is set to 0.98, the final temperature is set to 0.0001 and the number of iterations per fixed temperature is set to 50.

**Regular tabu search heuristic (Tabu):** The main idea behind a tabu search heuristic is to have a large number of iterations where in every iteration of the heuristic there is a single random pair exchange in the sequence. The heuristic is only allowed to make an exchange if this exchange improves the objective function and the exchange is not repeating a previous exchange that happened in the previous \( h \) iterations which has improved the objective function. The last \( h \) exchanges are kept in a list for checking. This list is called the tabu
list. This technique has been proven to be very useful for escaping isolation phenomena as well as escaping local optima during the course of search for a global optima. The following is an algorithmic description of the heuristic.

Regular Tabu Search Heuristic (Tabu)

Begin
    Initialize Tabu list with maximum size of 4
    Select a random sequence as current sequence
    Let \( T_1 = 0.1 \)
    While \( T_1 \geq 0.0001 \)
    Begin
        Repeat 50 times
        Begin
            Let \( L_1 \) = value of the objective function with current sequence
            Pick two random positions \( j \) and \( k \)
            If \( (j,k) \) is not in the Tabu list then
            Begin
                Swap jobs in the positions of \( j \) and \( k \)
                Let \( L_2 \) = value of the objective function after the swap
                If \( L_2 < L_1 \) then
                Begin
                    Add \( (j,k) \) to front of Tabu list
                    If Tabu maximum list size is exceeded, then
                    delete the item at the end of the list
                    Otherwise
                    Reverse and reject swap
                End If
            End If
        End Repeat
        Let \( T_1 = T_1 \times 0.98 \)
    End While
End Heuristic

For the proposed tabu search heuristic, setting the parameters is essential in achieving a good performance. After some experimentations, the parameters for the tabu search heuristic are set as follows; the total number of iterations is set to the same value of simulated annealing heuristic (for a fair comparison) and the tabu list size is set to four.

A new tabu search heuristic (Ntabu): The main idea behind the new tabu search heuristic is to introduce the concept of probability of accepting exchanges that do not improve the objective function into the tabu search heuristic. We were inspired to introduce this concept into a tabu search by observing the main concept behind the simulated annealing heuristic. This concept was integrated into the tabu search, which we call a new tabu search heuristic. The new tabu search heuristic is allowed to accept exchanges that are not in the tabu list and as in a regular tabu search, is an exchange that improves the objective function. It is also allowed in the new tabu search heuristic to have exchanges of a second type. The second type is an exchange that is not in the tabu list and does not improve the objective function with a small probability \( p \) that is reduced exponentially as the search for the global optima progresses. The following is an algorithmic description of the heuristic.

New Hybrid Tabu Search Heuristic (Ntabu)

Begin
    Initialize Tabu list with maximum size of 4
    Select a random sequence as current sequence
    Let \( T_1 = 0.1 \)
    While \( T_1 \geq 0.0001 \)
    Begin
        Repeat 50 times
        Begin
            Let \( L_1 \) = value of the objective function with current sequence
            Pick two random positions \( j \) and \( k \)
            If \( (j,k) \) is not in the Tabu list then
            Begin
                Swap jobs in the positions of \( j \) and \( k \)
                Let \( L_2 \) = value of the objective function after the swap
                Compute \( d \) and \( f \) where
                \[
                d = \frac{L_2 - L_1}{L_1}
                \]
                \[
                f = e^{-\frac{d}{T_1}}
                \]
                If \( (L_2 < L_1) \) or \( (L_2 > L_1 \text{ and with probability} f) \) then
                Begin
                    Add \( (j,k) \) to front of Tabu list
                    If Tabu maximum list size is exceeded, then
                    delete the item at the end of the list
                    Otherwise
                    Reverse and reject swap
                End If
            End If
        End Repeat
        Let \( T_1 = T_1 \times 0.98 \)
    End While
End Heuristic

For the proposed new tabu search heuristic, parameters are set to the same values as in simulated annealing and regular tabu search heuristics.

Experiments: According to the practitioner's experience, there is a roughly linear relation between the capacity of storage tank and the needed months for maintenance as shown in Table 1.

In addition, on the basis of the marketing experiences over 50 years, the predicted the maximum loads every month in one year as illustrated in Table 2.
Table 1: The capacity of tank and the needed months for maintenance

| Tank Capacity | 10 | 20 | 30 | 40 | 50 | 60 | 70 |
|---------------|----|----|----|----|----|----|----|
| Maintenance (month) | 2  | 3  | 4  | 5  | 6  | 7  | 8  |

Table 2: The maximal loads in a year

| Month | 1  | 2  | 3  | 4  | 5  | 6  |
|-------|----|----|----|----|----|----|
| Maximum Loads | 860 | 850 | 850 | 840 | 830 | 820 |
| Month | 7  | 8  | 9  | 10 | 11 | 12 |
| Maximum Loads | 830 | 820 | 810 | 850 | 830 | 840 |

Table 3: The attributes of 10 tanks

| Tank Number | 1  | 2  | 3  | 4  | 5  | 6  |
|-------------|----|----|----|----|----|----|
| Capacity (kiloliter) | 50 | 70 | 30 | 50 | 50 | 50 |
| Maintenance (months) | 6  | 8  | 4  | 6  | 6  | 6  |
| Tank number | 6  | 7  | 8  | 9  | 10 | 10 |
| Capacity (kiloliter) | 30 | 20 | 10 | 40 | 70 | 70 |
| Maintenance (months) | 4  | 3  | 2  | 5  | 8  | 8  |

Experimental design: In order to evaluate the performance of proposed approach, we adopt two kinds of data about storage tanks. First, we use the practical data of 10 tanks that are arranged to be maintained in certain year according to the dominating petroleum program. The capacities of these tanks and the corresponding month required are shown in Table 3. Second, the proposed algorithm is experimented on larger scale of problems by simulating 20- and 100 tank cases.

Three sets of data: one of them is the official data (10 tanks) from the dominating program in this year; the others (20 and 100 tanks) are generated randomly to simulate the performance in larger-scale problems.

Performance evaluation: To evaluate the performance of proposed approach, we use the relative error and standard deviation of the compared three heuristics SA, Tabu and Ntabu. Figure 2 and 3 illustrate the superiority of the Ntabu hyped evolutionary algorithm over the regular algorithms.

CONCLUSION

In this study, we introduced a new hybrid evolutionary algorithm and its implantation for solving real-world problem of oil storage tanks. The algorithm incorporates the American Petroleum Institute (API) standard 650 for open inspection procedure to produce a near optimal schedule for maintenance though continuously preserving a population diversity that ensures solution quality and convergence efficiency. The computational results show that the proposed evolutionary algorithm outperforms existing regular evolutionary algorithms and scheduler in literature and produces a higher quality solution which means better revenues though more operational tanks and more environmental protection though maximizing the utilization of existing tanks and less new tanks to build.

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