Quantum Capacities of Bosonic Channels

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We investigate the capacity of bosonic quantum channels for the transmission of quantum information. Achievable rates are determined from measurable moments of the channel by showing that every channel can asymptotically simulate a Gaussian channel which is characterized by second moments of the initial channel. We calculate the quantum capacity for a class of Gaussian channels, including channels describing optical fibers with photon losses, by proving that Gaussian encodings are optimal. Along the way we provide a complete characterization of degradable Gaussian channels and those arising from teleportation protocols.

One of the aims of Quantum Information Theory [1] is to follow the ideas of Shannon and to establish a theory of information based on the rules of quantum mechanics. A key problem along this way is the calculation of the quantum capacity of noisy quantum channels. That is, the question how much quantum information—measured in number of qubits—can be transmitted reliably per use of a given channel? Despite substantial progress [2] this can be answered only in very few cases [3] as a simple formula, comparable to Shannon’s coding theorem, is not known.

In this work we investigate the quantum capacity of bosonic channels which might describe transmission in space, as light sent through optical fibers, or in time, like in quantum memories [4]. The paper has two parts. In the first part we prove that the quantum capacity of any bosonic channel is lower bounded by that of a corresponding Gaussian channel $T_G$, which can be derived from measurable moments of $T$. This implies that for determining and certifying achievable rates for the transmission of quantum information through $T$ we need not know the channel exactly (which might be hardly possible in infinite dimensions), but merely its second moments, i.e., a few measurable parameters. In the second part we then explicitly calculate the quantum capacity of a class of Gaussian channels, which includes the important case of attenuation channels modelling optical fibers with photon losses and broad-band channels where losses and photon number constraints might be frequency dependent. Along the way we provide two tools that might be of independent interest: a complete characterization of degradable Gaussian channels and those arising from teleportation through Gaussian states.

PRELIMINARIES

Before we derive the main results we will briefly recall the basic notions [3, 5]. Consider a bosonic system of $N$ modes characterized by $N$ pairs of canonical operators $(Q_1, P_1, \ldots, Q_N, P_N) =: R$ for which the commutation relations $[R_k, R_l] = i\sigma_{kl}$ are governed by the symplectic matrix $\sigma$. The exponentials $W_\xi := e^{\xi R}$, $\xi \in \mathbb{R}^{2N}$ are called Weyl displacement operators. Their expectation value, the characteristic function, $\chi(\xi) := \text{tr}[\rho W_\xi]$ is the Fourier transform of the Wigner function and for Gaussian states it is a Gaussian

$$\chi(\xi) = e^{i\xi d - \frac{1}{2} \xi^T \Gamma \xi},$$

with first moments $d_k = \text{tr}[\rho R_k]$ and covariance matrix (CM) $\Gamma_{kl} := \text{tr}[\rho \{R_k - d_k, R_l - d_l\}]$. Note that coherent, squeezed and thermal states in quantum optics are all Gaussian states.

Gaussian channels [4, 7] transform Weyl operators as $W_\xi \mapsto W_X e^{-\frac{1}{2} Y Y^T} e^{1/2 Y^T \xi} Y$ and act on covariance matrices as

$$\gamma \mapsto X^T \gamma X + Y.$$  

Particularly important instances of single-mode Gaussian channels are attenuation and amplification channels for which $X = \sqrt{\eta} 1$ and $Y = |\eta - 1| 1$. For $0 \leq \eta \leq 1$ this models a single mode of an optical fiber with transmissivity $\eta$ where the environment is assumed to be in the vacuum state. The latter reflects the fact that thermal photons with optical frequencies are negligible at room temperature. For $\eta > 1$ the channel becomes an amplification channel, where the noise term $Y$ is now a consequence of the Heisenberg uncertainty.

Teleportation channels: We will now derive the form of Gaussian channels which are obtained when teleporting through a centered bipartite Gaussian state. As this is useful for applying but not necessary for understanding the following it might be skipped by the reader. Let $\Gamma = \begin{pmatrix} \Gamma_A & \Gamma_{AC} \\ \Gamma_{CA} & \Gamma_C \end{pmatrix}$ be the CM of a Gaussian state of $N_A + N_B$ modes with $N_A = N_B$. Assume Bob wants to teleport a quantum state of $N_B$ modes with CM $\gamma$ to Alice. Using the standard protocol [8] he sends pairs of modes from $\gamma$ and $\Gamma_B$ through 50:50 beam-splitters, measures the $Q$ and $P$ quadratures, and then communicates the outcomes. Depending on the latter Alice applies displacements to the modes in $\Gamma_A$. The simplest way of
deriving an expression for the output is to start with the Wigner representation and to assume that the state to be teleported is a centered Gaussian. The Wigner function before the measurement is up to normalization given by \( \exp -\xi [M_{\text{BS}}^T (\Gamma \otimes \gamma) ^{-1} M_{\text{BS}}] \xi \), where \( M_{\text{BS}} \) corresponds to the beam-splitter operation. With \( \xi = (\xi_A, \xi_B, \xi_{B'}) \) the final Wigner function is then proportional to

\[
\int d\xi_B d\xi_{B'} e^{-\xi [M_X^T M_{\text{BS}}^T (\Gamma \otimes \gamma) ^{-1} M_{\text{BS}} M_X] \xi},
\]

where \( M_X \) incorporates the displacements, i.e., it is the identity matrix plus an arbitrary \( 2N_B \times 2N_B \) off-diagonal block which maps the \( 2N_B \) measurement outcomes onto the respective displacements. In order to circumvent integrating Eq. 3 we can now go to the characteristic function, i.e., the Fourier transformed picture. The integration then boils down to picking out the upper left block of the inverted matrix \( [M_X^T M_{\text{BS}}^T (\Gamma \otimes \gamma) ^{-1} M_{\text{BS}} M_X] ^{-1}. \)

The inversion is, however, trivial since \( M_{\text{BS}}^{-1} = M_{\text{BS}}^T \) and \( M_X^{-1} \) is obtained from \( M_X \) by changing the sign of all off-diagonal entries. In this way we obtain that the input CM is transformed to

\[
\gamma \mapsto X^T \gamma X + \left[ \Gamma_A + \Gamma C \Lambda X + (\Gamma C \Lambda X)^T \gamma^T \Lambda \gamma X \right],
\]

where \( \Lambda = \text{diag}(1, -1, 1, -1, \ldots) \) and \( X \) is such that \( \sqrt{2}X \) is the matrix of displacement transformations, i.e., the gain which is typically chosen to be \( \sqrt{2} \).

Clearly, Eq. 4 has the form \( \chi_{\text{in}} \) and following the above lines it is straightforward to show that the channel is Gaussian and maps any (not necessarily centered Gaussian) input characteristic function \( \chi_{\text{in}} \) into

\[
\chi_{\text{out}}(\xi) = \chi_{\text{in}}(X \xi) \chi_{\text{in}}(\xi \oplus \Lambda X \xi).
\]

For standard protocols \( (X = I) \) on single modes \( (N_A = N_B = 1) \) this was derived in [3].

**ACHIEVABLE RATES FOR ARBITRARY CHANNELS**

The subject of interest is the quantum capacity \( Q(T) \) of an arbitrary—a priori unknown—channel \( T \). We will show how one can certify achievable rates for the transmission of quantum information through \( T \) by only looking at the CM \( \Gamma \) of a state \( \rho_T = (T \otimes \text{id}) (\psi) \) which is obtained by sending half of an arbitrary entangled state \( \psi \) through the channel. \( \Gamma \) could be determined by homodyne measurements. The argument combines (i) the relation between entanglement distillation and quantum capacities observed in [10], (ii) the extremality of Gaussian states shown in [11] and (iii) the explicit form of Gaussian teleportation channels derived in the previous section. All together this leads to the chain of inequalities

\[
Q(T) \geq D_{\rightarrow} (\rho_T) \geq D_{\rightarrow} (\mathcal{G}(\rho_T)) \geq Q(T_G).
\]

Here \( D_{\rightarrow} (\rho_T) \) is the distillable entanglement under protocols with one-way communication (from Bob to Alice). Since a classical side channel does not increase \( Q(T) \) this is clearly a lower bound to the capacity as Alice and Bob could simply first distill \( \rho_T \) and then use the obtained maximally entangled states for teleportation [11]. The second inequality uses that replacing \( \rho_T \) by a Gaussian state \( \mathcal{G}(\rho_T) \) with the same CM \( \Gamma \) can only decrease the distillable entanglement [11]. Finally, if we use the Gaussian state in turn as a resource for establishing a teleportation channel \( T_G \) we end up with the sought inequality \( Q(T) \geq Q(T_G) \). \( T_G \) is then the Gaussian channel in Eqs. [12, 13], which is for a fixed teleportation protocol (a fixed matrix \( X \)) completely determined by \( \Gamma \).

Bounds on the quantum capacity of Gaussian channels were derived in [12, 13] and we will show below that it can be calculated exactly for some important cases. Note that a simple bound for \( Q(T) \) can be obtained from a lower bound to \( D_{\rightarrow} (\mathcal{G}(\rho_T)) \), the conditional entropy of the Gaussian state with CM \( \Gamma \), i.e., \( Q(T) \geq S(\Gamma_A) - S(\Gamma) \).

Before we proceed, two comments on the quality of the above bound and its operational meaning are in order: The given argument holds for arbitrary \( T \) and \( \psi \). However, since we bound by Gaussian quantities the inequality might become trivial (i.e., \( Q(T_G) = 0 \) though \( Q(T) \gg 0 \)) if both \( T \) and \( \psi \) are too far from being Gaussian. On the other hand, if \( T \) is Gaussian and \( |\psi\rangle = (\cosh r \gamma)^{-1} \sum_n (\tanh r)^n |nn\rangle \) is a two-mode squeezed state, then in the limit \( r \to \infty \) the inequality becomes tight, i.e., \( Q(T_G) \to Q(T) \) with exponentially vanishing gap.
tional meaning: sender and receiver first establish some entanglement, distill it and then use it as a resource for teleportation. The second step can also be understood in operational terms. To achieve $\rho_T \rightarrow \mathcal{G}(\rho_T)$ both sender and receiver have to apply an array of beam-splitters (see Fig.1) to many copies of the shared state $\rho_T$. Asymptotically every reduced state at the output will then converge to $\mathcal{G}(\rho_S)$ [11]. By applying this 

**QUANTUM CAPACITY OF GAUSSIAN CHANNELS**

It was proven in [2] that the quantum capacity of a quantum channel $T$ can be expressed as

$$Q(T) = \lim_{n \to \infty} \frac{1}{n} \sup_{\rho} \, J(\rho, T^\otimes n) ,$$

(7)

$$J(\rho, T) = S(T(\rho)) - S((T \otimes \text{id})(\psi)) ,$$

(8)

where $\psi$ is a purification of $\rho$ and $J$ is known as the coherent information. In general, the calculation of $Q(T)$ from the above formula is a daunting task since (i) the coherent information is known to be not additive, i.e., the regularisation $n \to \infty$ is necessary, and (ii) due to lacking concavity properties there are local maxima which are not global ones. On top of this, for bosonic channels the optimization is over an an infinite dimensional space.

Fortunately, for a class of Gaussian channels including the important case of the lossy channel, these obstacles can be circumvented by exploiting recent results on degradability of channels [3,11] and extremality of Gaussian states [11].

To this end consider a channel $T(\rho_S) = \text{tr}_E[U(\rho_S \otimes \varphi_E)U^\dagger]$ expressed in terms of a unitary coupling between the system $S$ and the environment $E$ which is initially in a pure state $\varphi_E$. The conjugate channel $T_c(\rho_S) = \text{tr}_E[U(\rho_S \otimes \varphi_E)U^\dagger]$ is defined as a mapping from the system to the environment. As shown in [3] the coherent information can be expressed in terms of a conditional entropy if there exists a channel $T'$ such that $T' \circ T = T_c$ — in this case $T$ is called degradable. More precisely, if $\hat{\rho}_S|E'\rangle\langle E'|$ is the extension of the state $\hat{\rho}_S = T' \circ T(\rho)$ to the environment $E'$ of $T'$, then

$$J(\rho, T) = S(\hat{\rho}_S|E') - S(\hat{\rho}_S) =: S(E'|S') .$$

(9)

The conditional entropy $S(E'|S')$ is known to be strongly sub-additive [1], i.e., for a composite system $S(E'|S'_{12}) \leq S(E'|S_1) + S(E'|S_2)$. This has important consequences: for a set $\{T_i\}$ of degradable channels $J(\rho, \otimes_i T_i) \leq \sum_i J(\rho_i, T_i)$, where $\rho_i$ are the corresponding reduced states, and if each $T_i$ is a Gaussian channel, we have in addition

$$J(\rho, \otimes_i T_i) \leq \sum_i J(\rho_i, T_i) \leq \sum_i J(\mathcal{G}(\rho_i), T_i) .$$

(10)

The last inequality follows from the extremality of Gaussian states w.r.t. the conditional entropy [3,11] together with the fact that for Gaussian channels $T_c$ can be chosen to be Gaussian and the CM is transformed irrespective of whether the input was Gaussian or not. As a consequence, if $T_i$ are degradable Gaussian channels, then

$$Q(\otimes_i T_i) = \sum_i \sup_{\rho_G} J(\rho_G, T_i) ,$$

(11)

where the supremum is now taken only over Gaussian input states $\rho_G$. Calculating the latter for Gaussian channels is now a feasible task which was solved for the single-mode case in [12] and in [13] for broadband channels under power constraints using Lagrange multipliers. In fact, if we impose a constraint on the input energy of the form $\sum_i \omega_i N_i = E$, where $N_i$ is the average input photon number of mode $i$ with corresponding frequency $\omega_i$, then the above argumentation still holds, since the constraint just depends on the CM. The importance of Eq.(11) stems from the fact that a large class of Gaussian channels is indeed degradable, as shown in [14] and extended below. In particular we can apply Eq.(11) to attenuation (amplification) channels with transmissivity $\eta$ (gain $\sqrt{\eta}$). Together with the optimization carried out in [12] this yields $Q$ as a function of $\eta$ (see Fig.2):

$$Q(\eta) = \max \left\{ 0, \log_2 |\eta| - \log_2 (1 - |\eta|) \right\} .$$

(12)

Note that the quantum capacity of every degradable Gaussian channel can easily be calculated as $J$ becomes a concave function of the CM such that local maxima are
DEGRADABLE GAUSSIAN CHANNELS

We will now investigate the condition under which Eq. (11) was derived and characterize the set of degradable Gaussian channels—extending the results of [14]. To this end we represent the channel in terms of a unitary normal form [17] which only depends on the symplectic invariant det $X$ such that $Q(T)$ of every such channel is given by Eq. (12) with $\eta = \det X$.

Let us finally briefly comment on the case $N_E \neq N_S$. If the environment is smaller than the system, then we can easily follow the above lines for instance by choosing a representation of the channel with larger $N_E$ equal to $N_S$ [18]. It is worth mentioning that if $S$ corresponds to a passive (i.e., number preserving) operation, then for $N_E < N_S$ there are always unaffected modes such that $Q(T) = \infty$ without additional constraints. If $N_E > N_S$ then Eq. (11) is merely a necessary, whereas Eq. (14) is still a necessary and sufficient condition for degradability [18]. Applying the latter to a general single-mode channel with $N_S = 1, N_E = 2$ shows that generically one has neither degradability nor anti-degradability.

Hence, it remains open whether in this case the capacity is given by Eq. (11). However, we can easily derive an upper bound by exploiting the fact that every Gaussian channel $T$ can be decomposed as $T = T_1 \circ T_2$, where $T_2$ is a minimal noise channel [4] for which $N_E = N_S$ with $X_2 = X$, $Y_2 \leq Y$ and $T_1$ is a classical noise channel for which $X_1 = 1$, $Y_1 = Y - Y_2$. Due to the bottleneck-inequality for capacities (cf. [12]) we have $Q(T) \leq Q(T_2)$ where the latter is in the single-mode case again given by Eq. (12) with $\eta = \det X$. A lower bound is always given by the r.h.s. of Eq. (11) as calculated in [12].

In summary we characterized the set of degradable Gaussian channels and showed that their quantum capacity can be calculated as it is attained for Gaussian product inputs. For arbitrary non-Gaussian channels we derived a certifiable Gaussian lower bound. Both ideas can be applied to finite dimensional systems as well. This will, however, be the content of future work [19].

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The symplectic condition $S \sigma S^T = \sigma$ amounts to $A \sigma A^T + B \sigma B^T = C \sigma C^T + D \sigma D^T = \sigma$ and $A \sigma C^T + B \sigma D^T = 0$. To see this one can exploit the fact that $\frac{1}{2} (\mathbb{1} \pm i \sigma)$ are orthogonal projectors. In this way Eq. (14) can be expressed in terms of the two conditions $[\sigma, K] = 0$ and $\mathbb{1} \geq KK^T$ which can in turn be rephrased as
\[
(X \sigma X^T \sigma Y) - (X \sigma X^T \sigma Y)^T = 0, \quad (16)
\]
\[
(X \sigma X^T \sigma Y) + (X \sigma X^T \sigma Y)^T + Y \leq 0, \quad (17)
\]
which together imply Eq. (15).

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