Extensions of the Minimum Cost Homomorphism Problem

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Abstract. Assume $D$ is a finite set and $R$ is a finite set of functions from $D$ to the natural numbers. An instance of the minimum $R$-cost homomorphism problem ($\text{MinHom}_R$) is a set of variables $V$ subject to specified constraints together with a positive weight $c_{vr}$ for each combination of $v \in V$ and $r \in R$. The aim is to find a function $f : V \rightarrow D$ such that $f$ satisfies all constraints and $\sum_{v \in V} \sum_{r \in R} c_{vr} r(f(v))$ is minimized. This problem unifies well-known optimization problems such as the minimum cost homomorphism problem and the maximum solution problem, and this makes it a computationally interesting fragment of the valued CSP framework for optimization problems. We parameterize $\text{MinHom}_R(\Gamma)$ by constraint languages, i.e., sets $\Gamma$ of relations that are allowed in constraints. A constraint language is called conservative if every unary relation is a member of it; such constraint languages play an important role in understanding the structure of constraint problems. The dichotomy conjecture for $\text{MinHom}_R(\Gamma)$ is the following statement: if $\Gamma$ is a constraint language, then $\text{MinHom}_R(\Gamma)$ is either polynomial-time solvable or NP-complete. For $\text{MinHom}$ the dichotomy result has been recently obtained [Takhanov, STACS, 2010] and the goal of this paper is to expand this result to the case of $\text{MinHom}_R$ with conservative constraint language. For arbitrary $R$ this problem is still open, but assuming certain restrictions on $R$ we prove a dichotomy. As a consequence of this result we obtain a dichotomy for the conservative maximum solution problem.

1 Introduction

Constraint satisfaction problems (CSP) and valued constraint satisfaction problems (VCSP) are natural ways of formalizing a large number of computational problems arising in combinatorial optimization, artificial intelligence, and database theory. CSP has the following two equivalent formulations: (1) to find an assignment of values to a given set of variables, subject to constraints on the values that can be assigned simultaneously to specified subsets of variables, and (2) to find a homomorphism between two finite relational structures $A$ and $B$. VCSP is a “soft” version of CSP where constraint relations are replaced by functions from set of tuples to some totally ordered set with addition operation (for example, rational numbers). A solution is defined as an assignment
to variables that maximize a functional which is equal to a sum of constraint functions applied to corresponding variables. Applications of CSPs and VCSPs arise in the propositional logic, database and graph theory, scheduling, biology and many other areas. CSP and its subproblems has been intensively studied by computer scientists and mathematicians since the 70s, and recently attention has been paid to its modifications such as VCSP. Considerable attention has been given to the case where the constraints are restricted to a given finite set of relations \( \Gamma \), called a constraint language [2, 5, 12, 20]. For example, when \( \Gamma \) is a constraint language over the boolean set \( \{0, 1\} \) with four ternary predicates \( x \lor y \lor z, x \lor y \lor z, x \lor y \lor z, x \lor y \lor z \) we obtain 3-SAT. For every constraint language \( \Gamma \), it has been conjectured that CSP(\( \Gamma \)) is either in P or NP-complete [5]. An analogous situation appears in VCSP where the constraint language is defined as a set of “soft” predicates.

We believe that problems like minimum cost homomorphism problem (MinHom) has an intermediate position between CSPs and VCSPs which makes their structure important for understanding the relationship between “hard” and “soft” constraints in optimization. In the minimum cost homomorphism problem, we are given variables subject to constraints and, additionally, costs on variable/value pairs. Now, the task is not just to find any satisfying assignment to the variables, but one that minimizes the total cost. In the context of VCSP this is equivalent to addition of “soft” constraints equal to characteristic functions of one element sets. We will consider a weighted version of this problem.

**Definition 1.** Suppose we are given a finite domain set \( A \), a finite constraint language \( \Gamma \subseteq \bigcup_{k=1}^{\infty} 2^{A^k} \) and a finite set of functions \( R \subseteq \{ r : A \rightarrow \mathbb{N} \} \). Denote by MinHom\(_R\)(\( \Gamma \)) the following minimization task:

**Instance:** A triple \((V, C, W)\) where

- \( V \) is a set of variables;
- \( C \) is a set of constraints, where each constraint \( C \in C \) is a pair \((s, \rho)\), such that
  - \( s = (v_1, \ldots, v_m) \) is a tuple of variables of length \( m \), called the constraint scope;
  - \( \rho \) is an element of \( \Gamma \) with arity \( m \), called the constraint relation.
- Weights \( w_{vr} \in \mathbb{N}, v \in V, r \in R \).

**Solution:** A function \( f \) from \( V \) to \( A \), such that, for each variable \( v \in V, f(v) \in A \), and for each constraint \((s, \rho) \in C \), with \( s = (v_1, \ldots, v_m) \), the tuple \((f(v_1), \ldots, f(v_m))\) belongs to \( \rho \).

**Measure:** \( \sum_{v \in V} \sum_{r \in R} w_{vr}r(f(v)). \)

**Definition 2.** For \( R^* = \{ e_i | i \in A \} \), MinHom\(_{R^*}\)(\( \Gamma \)) is called minimum cost homomorphism problem where \( e_i : A \rightarrow \mathbb{N} \) denotes a characteristic function of \( \{i\} \subseteq A \).
We will write $\text{MinHom}$ instead of $\text{MinHom}_R$ for short. $\text{MinHom}$ has applications in defence logistics [10] and machine learning [4]. Complete classification of constraint languages $\Gamma$ for which $\text{MinHom}(\Gamma)$ is polynomial-time solvable has recently been obtained in [21]. The question for which directed graphs $H$ the problem $\text{MinHom}(\{H\})$ is polynomial-time solvable was considered in [7–11]. Maximum Solution Problem ($\text{MaxSol}$), which is defined analogously to $\text{MinHom}$, but with $A \subseteq \mathbb{N}$ and a functional of the form $\sum_{v \in V} w_v f(v)$ to maximize, was investigated in a series of papers [14–16]. It is easy to see that if $n = \max s + 1$ and $R = \{n - x\}$, then $\text{MinHom}_R(\Gamma) = \text{MaxSol}(\Gamma)$.

In this paper, we will assume that a constraint language $\Gamma$ contains all unary predicates over a domain set $A$ and approach the problem of characterizing the complexity of $\text{MinHom}_R(\Gamma)$ in its most general form by algebraic methods. When $R$ satisfies certain conditions, we obtain a dichotomy for $\text{MinHom}_R(\Gamma)$, i.e., if $\text{MinHom}_R(\Gamma)$ is not polynomial-time solvable, then it is NP-hard. As a consequence, we obtain a dichotomy for conservative $\text{MaxSol}$.

In Section 2, we present some preliminaries together with results connecting the complexity of $\text{MinHom}_R$ with conservative algebras. The main dichotomy theorem is stated in Section 3 and its proof is divided into several parts which can be found in Sections 4-6. Finally, in Section 7 we present directions for future research.

2 Algebraic structure of tractable constraint languages

Recall that an optimization problem $A$ is called NP-hard if some NP-complete language can be recognized in polynomial time with the aid of an oracle for $A$. We assume that $P \neq \text{NP}$.

**Definition 3.** Suppose we are given a finite set $A$ and a constraint language $\Gamma \subseteq \bigcup_{k=1}^{\infty} 2^{A^k}$. The language $\Gamma$ is said to be $R$-tractable if, for every finite subset $\Gamma' \subseteq \Gamma$, the task $\text{MinHom}_R(\Gamma')$ is polynomial-time solvable, and $\Gamma$ is called $R$-NP-hard if there is a finite subset $\Gamma' \subseteq \Gamma$, such that the task $\text{MinHom}_R(\Gamma')$ is NP-hard.

First, we will state some standard definitions from universal algebra.

**Definition 4.** Let $\rho \subseteq A^n$ and $f : A^n \rightarrow A$. We say that the function (operation) $f$ preserves the predicate $\rho$ if, for every $(x_1^i, \ldots, x_m^i) \in \rho, 1 \leq i \leq n$, we have that $(f(x_1^i, \ldots, x_m^i), \ldots, f(x_1^i, \ldots, x_m^i)) \in \rho$.

For a constraint language $\Gamma$, let $\text{Pol}(\Gamma)$ denote the set of operations preserving all predicates in $\Gamma$. Throughout the paper, we let $A$ denote a finite domain and $\Gamma$ a constraint language over $A$. We assume the domain $A$ to be finite.

**Definition 5.** A constraint language $\Gamma$ is called a relational clone if it contains every predicate expressible by a first-order formula involving only
a) predicates from $\Gamma \cup \{=^A\}$;  
b) conjunction; and  
c) existential quantification.

First-order formulas involving only conjunction and existential quantification are often called primitive positive (pp) formulas. For a given constraint language $\Gamma$, the set of all predicates that can be described by pp-formulas over $\Gamma$ is called the closure of $\Gamma$ and is denoted by $\langle \Gamma \rangle$.

For a set of operations $F$ on $A$, let $\text{Inv}(F)$ denote the set of predicates preserved under the operations of $F$. Obviously, $\text{Inv}(F)$ is a relational clone.

The next result is well-known [1, 6].

**Theorem 1.** For a constraint language $\Gamma$ over a finite set $A$, $\langle \Gamma \rangle = \text{Inv}(\text{Pol}(\Gamma))$.

Theorem 1 tells us that the Galois closure of a constraint language $\Gamma$ is equal to the set of all predicates that can be obtained via pp-formulas from the predicates in $\Gamma$. We will omit the proof of the following standard theorem.

**Theorem 2.** For any finite constraint language $\Gamma$, finite $\Gamma' \subseteq \langle \Gamma \rangle$ and finite $R \subseteq \{r : A \rightarrow \mathbb{N}\}$ there is a polynomial time reduction from $\text{MinHom}_R(\Gamma')$ to $\text{MinHom}_R(\Gamma)$.

The previous theorem tells us that the complexity of $\text{MinHom}_R(\Gamma)$ is basically determined by $\text{Inv}(\text{Pol}(\Gamma))$, i.e., by $\text{Pol}(\Gamma)$. That is why we will be concerned with the classification of sets of operations $F$ for which $\text{Inv}(F)$ is a tractable constraint language.

**Definition 6.** An algebra is an ordered pair $A = (A, F)$ such that $A$ is a nonempty set (called a universe) and $F$ is a family of finitary operations on $A$. An algebra with a finite universe is referred to as a finite algebra.

**Definition 7.** An algebra $A = (A, F)$ is called $R$-tractable if $\text{Inv}(F)$ is a $R$-tractable constraint language and $A$ is called $R$-NP-hard if $\text{Inv}(F)$ is an $R$-NP-hard constraint language.

For $B \subseteq A$, define $R_B = \{f|_B | f \in R\}$, where $f|_B$ is a restriction of $f$ on a set $B$. We will use the term MinHom-tractable (NP-hard) instead of $\{e_i | i \in A\}$-tractable (NP-hard) and, in case $A = \{0, 1\}$, the term min-tractable (NP-hard) instead of $\{x\}$-tractable (NP-hard).

We only need to consider a very special type of algebras, so called conservative algebras.

**Definition 8.** An algebra $A = (A, F)$ is called conservative if for every operation $f \in F$ we have that $f(x_1, \ldots, x_n) \in \{x_1, \ldots, x_n\}$.

Since we assume that $\Gamma$ is a constraint language with all unary relations over the domain set $A$, then $A = (A, \text{Pol}(\Gamma))$ is conservative. Besides conservativeness of constraint languages we will make some additional restrictions on function sets $R$. 

Definition 9. Suppose we are given a finite set of functions \( R \subseteq \{ r : A \to \mathbb{N} \} \). Denote by \( G(R) = (A,E(R)) \) a directed graph with a set of vertices \( A \) and an edge set \( E(R) = \{(a,b) | \exists r \in R r(a) > r(b) \} \). The \( UG(R) \) is the graph \( G(R) \) with all edges considered as undirected. We will call \( G \) a preference graph and \( UG \) an undirected preference graph.

In the sequel, we will assume that a graph \( UG \) is complete. It is easy to see that \( UG(\{e_i | i \in A\}) \), \( UG(\{n-x\}) \) are complete and our results can be applied to \( MinHom \) and \( MaxSol \).

3 Boolean case and the necessary local conditions

The first step to understand the structure of \( R \)-tractable algebras is to understand the boolean case. Well-known structure of boolean clones [18] helps us to prove the following theorem.

Theorem 3. A boolean clone \( H \) is \( MinHom \)-tractable if either \( \{x \land y, x \lor y\} \subseteq H \) or \( \{(x \land \overline{y}) \lor (\overline{y} \land z) \lor (x \land z)\} \subseteq H \), where \( \overline{x} \), \( x \land y \), \( x \lor y \) denote negation, conjunction and disjunction. Otherwise, \( H \) is \( MinHom \)-NP-hard. A conservative boolean clone \( H \) is min-tractable if either \( \{x \land y\} \subseteq H \) or \( \{(x \land \overline{y}) \lor (\overline{y} \land z) \lor (x \land z)\} \subseteq H \). Otherwise, \( H \) is min-NP-hard.

In the proof of Theorem 3 we will need the following definition.

Definition 10. A constraint language \( \Gamma \) over \( \{0,1\} \) is called a \( MinHom(min) \)-maximal constraint language if it is conservative \( MinHom(min) \)-tractable and is not contained in any other conservative \( MinHom(min) \)-tractable languages.

We identify all \( MinHom(min) \)-maximal constraint languages using Post’s classification [18]. Via Theorems 1,2 we conclude that every \( MinHom(min) \)-maximal constraint language corresponds to some conservative functional clone. In the case \( A = \{0,1\} \), there is a countable number of conservative clones: we list them below according to the table on page 76 [17].

Lemma 1. The relational clones \( Inv(M_{01}) \) and \( Inv(S_{01}) \) are \( MinHom \)-maximal constraint languages. Every other constraint language given in the table that is not contained in any of these two is \( MinHom \)-NP-hard. The relational clones \( Inv(K_{01}) \) and \( Inv(S_{01}) \) are min-maximal constraint languages. Every other constraint language given in the table that is not contained in any of these two is min-NP-hard.
Proof. For every row, the closure of the predicates given is equal to the set of all predicates preserved under the functions of the corresponding clone.

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| $T_{01}$ | $x = 0, x = 1$ |
| $M_{01}$ | $x = 0, x = 1, x_1 \leq x_2$ |
| $S_{01}$ | $x = 0, x_1 \neq x_2$ |
| $SM$ | $x_1 \neq x_2, x_1 \leq x_2$ |
| $L_{01}$ | $x = 1, x_1 \oplus x_2 \oplus x_3 = 0$ |
| $U_{01}$ | $x = 0, x = 1, x_1 = x_2 \lor x_1 = x_3$ |
| $K_{01}$ | $x = 0, x = 1, x_1 = x_2 x_3$ |
| $D_{01}$ | $x = 0, x = 1, x_1 = x_2 \lor x_3$ |
| $I_{m}^{n}$ | $x = 1, x_1 x_2 \ldots x_m = 0$ |
| $MI_{I}^{n}$ | $x_1, x_2 \leq x_2, x_1 x_2 \ldots x_m = 0$ |
| $O_{m}^{n}$ | $x = 0, x_1 \lor x_2 \lor \ldots \lor x_m = 1$ |
| $MO_{m}^{n}$ | $x = 0, x_1 \leq x_2, x_1 \lor x_2 \lor \ldots \lor x_m = 1$ |

where $x \oplus y = x + y (mod \ 2)$.

The class $Inv(T_{01})$ is MinHom(min)-tractable, since it contains only two simple unary predicates $\{0\}$ and $\{1\}$. As we will see later, it cannot be MinHom(min)-maximal since it is included in other MinHom(min)-tractable constraint languages.

Let us prove that $Inv(M_{01})$ and $Inv(S_{01})$ are MinHom-tractable. By Theorem 2, it is equivalent to polynomial solvability of $MinHom(\{\{0\} , \{1\} , \{(x_1, x_2) | x_1 \leq x_2\})$ and $MinHom(\{\{0\} , \{(x_1, x_2) | x_1 \neq x_2\})$, because the classes $Inv(M_{01})$ and $Inv(S_{01})$ are the closures of those sets. We will skip the proof since it can be found in [13] (in this paper boolean MinHom called Max AW Ones). It is easy to see that MinHom-tractability implies min-tractability of those classes.

Let us show that all the classes in the table, except $Inv(M_{01})$, $Inv(S_{01})$, and $Inv(T_{01})$, are MinHom-NP-hard, and all the classes, except $Inv(M_{01})$, $Inv(S_{01})$, $Inv(T_{01})$, $Inv(K_{01})$ and $Inv(MI_{I}^{n})$, $\ldots$, $Inv(MI_{m}^{n})$, where $MI_{I}^{n} = \bigcup_{m=1}^{\infty} MI_{m}^{n}$, are min-NP-hard. Since,}

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| $x_1 \lor x_2 = \exists x_3 [x_1 \neq x_3] \land [x_3 \leq x_2]$ |
| $x_1 \lor x_2 = \exists x_3 [x_3 = 1] \land [x_3 = x_1 \lor x_3 = x_2]$ |
| $\bar{x_1} \lor \bar{x_2} = \exists x_3 [x_3 = 0] \land [x_3 = x_1 x_2]$ |
| $\bar{x_1} \lor \bar{x_2} = \exists x_3 [x_3 = 1] \land [x_3 = x_1 \lor x_2]$ |
| $\bar{x_1} \lor \bar{x_2} = \exists x_3 \ldots x_m [x_1 x_2 \ldots x_m = 0] \land [x_2 = x_3] \land \ldots \land [x_m = x_m] |
| $x_1 \lor x_2 = \exists x_3 \ldots x_m [x_1 \lor x_2 \ldots \lor x_m = 1] \land [x_2 = x_3] \land \ldots \land [x_m = x_m]$ |

we see that \{(x_1, x_2) | x_1 \lor x_2 \} $\in Inv(SM)$, $Inv(U_{01})$, $Inv(D_{01})$, $Inv(O_{m}^{n})$, $Inv(MO_{m}^{n})$ and \{(x_1, x_2) | \bar{x_1} \lor \bar{x_2} \} $\in Inv(K_{01})$, $Inv(I_{m}^{n})$, $Inv(MI_{m}^{n})$.

We first prove that $MinHom_{\{x\}}(\{(x_1, x_2) | x_1 \lor x_2 \})$ is NP-hard. Suppose an instance of this problem consists of an undirected graph $G = (V,E)$ where each vertex is considered as a variable. For each pair of variables
Proof of Theorem 12. Obviously, \( \{\wedge, \vee\} \) is a basis of \( M_{01} \), \( \{\wedge\} \) is a basis of \( K_{01} \) and \( \{x \wedge \overline{y} \} \lor (\overline{y} \wedge z) \lor (x \wedge z) \) is a basis of \( S_{01} \). By Lemma 1 we conclude the statement of the theorem. \( \square \)

**Definition 11.** We call a pair of vertices \( \{a, b\} \) in \( G(R) \) a MinHom-pair if they have arcs in both directions and we call \( \{a, b\} \) a min-pair if they have an arc in one direction only.

Suppose \( f \in F \). By \( \frac{a}{b} f \), we mean \( a \neq b \) and \( f(a, b) = f(b, a) = b \). Guided by Theorem 3 we obtain the following definition.
Definition 12. Let $F$ be a conservative functional clone over $A$ and $R \subseteq \{ r : A \to \mathbb{N} \}$. We say that $F$ satisfies the necessary local conditions for $R$ if and only if

1. for every MinHom-pair $\{ a, b \}$ of $G(R)$, either
   1.a) there exists $f_1, f_2 \in F$ s.t. $\frac{a}{b} f_1$ and $\frac{a}{b} f_2$; or
   1.b) there exists $f \in F$ s.t. $f|_{\{a,b\}}(x,x,y) = f|_{\{a,b\}}(y,x,x) = f|_{\{a,b\}}(y,x,y) = y$.

2. for every min-pair $\{ a, b \}$ of $G(R)$ such that $(a, b) \in E(R)$, either
   2.a) there exists $f \in F$ s.t. $\frac{a}{b} f$; or
   2.b) there exists $f \in F$ s.t. $f|_{\{a,b\}}(x,x,y) = f|_{\{a,b\}}(y,x,x) = f|_{\{a,b\}}(y,x,y) = y$.

Theorem 4. Suppose $F$ is a conservative functional clone. If $F$ is $R$-tractable and $UG(R)$ is complete, then it satisfies the necessary local conditions for $R$. If $F$ does not satisfy the necessary local conditions for $R$, then it is $R$-NP-hard.

As in case of $MinHom$, the necessary local conditions are not sufficient for $R$-tractability of a conservative clone. Let $M = \{ B \mid B \subseteq A, |B| = 2, F|_B \text{ contains 2 different binary commutative functions} \}$, $M^o = \{ (a, b) \mid (a, b) \in M \}$ and $M_c = \{ B \mid B \subseteq A, |B| = 2 \} \setminus M$.

Introduce an undirected graph without loops $T^R_B = (M^o \cap E(R), P)$ where $P = \{ \langle (a, b), (c, d) \rangle \mid (a, b), (c, d) \in M^o \cap E(R) \}$, there is no $f \in F : \frac{a}{b} \frac{c}{d} f$.

Theorem 5. Suppose $F$ satisfy the necessary local conditions for $R$ and $UG(R)$ is complete. If the graph $T^R_B = (M^o \cap E(R), P)$ is bipartite, then $F$ is $R$-tractable. Otherwise, $F$ is $R$-NP-hard.

A proof for NP-hard case of Theorem 5 will be omitted since it is basically the same as in case of $MinHom$ [21].

4 Multi-sorted MinHom and its tractable case

Definition 13. For any collection of sets $A = \{ A_i \mid i \in I \}$, and any list of indices $i_1, \ldots, i_m \in I$, a subset $\rho$ of $A_{i_1} \times \ldots \times A_{i_m}$, together with the list $(i_1, \ldots, i_m)$, will be called a multi-sorted relation over $A$ with arity $m$ and signature $(i_1, \ldots, i_m)$. For any such relation $\rho$, the signature of $\rho$ will be denoted $\sigma(\rho)$.

Definition 14. Let $\Gamma$ be a set of multi-sorted relations over a collection of sets $A = \{ A_i \mid i \in I \}$. The multi-sorted $MinHom$ problem over $\Gamma$, denoted $MMinHom(\Gamma)$, is defined to be the minimization problem with

Instance: A quadruple $(V, \delta, C, W)$ where
- $V$ is a set of variables;
- $\delta$ is a mapping from $V$ to $I$, called the domain function;
Definition 15. A set of multi-sorted relations over $M_{\text{MinHom}}$.

Definition 16. Let $\Gamma$ be a collection of sets. An $n$-ary multi-sorted operation $t$ on $A$ is defined by a collection of interpretations $\{t^A | A \in A\}$, where each $t^A$ is an $n$-ary operation on the corresponding set $A$. The multi-sorted operation $t$ on $A$ is said to be a polymorphism of a multi-sorted relation $\rho$ over $A$ with signature $(\delta(1), \ldots, \delta(m))$ if, for any $(a_{11}, \ldots, a_{1n}), \ldots, (a_{mn}, \ldots, a_{mn}) \in \rho$, we have

$$
t \left( \begin{array}{c} a_{11} \cdots a_{1n} \\ \vdots \\ a_{m1} \cdots a_{mn} \end{array} \right) = \left( \begin{array}{c} t^{\delta(1)}(a_{11}, \ldots, a_{1n}) \\ \vdots \\ t^{\delta(m)}(a_{m1}, \ldots, a_{mn}) \end{array} \right) \in \rho$$

For any given set of multi-sorted relations $\Gamma$, $M\text{Pol}(\Gamma)$ denotes the set of multi-sorted operations which are polymorphisms of every relation in $\Gamma$.

Definition 17. Suppose a set of operations $H$ over $D$ is conservative and $B \subseteq \{\{x, y\} | x, y \in D, x \neq y\}$. A pair of binary operations $\phi, \psi \in H$ is called a tournament pair on $B$, if $\forall \{x, y\} \in B \phi(x, y) = \phi(y, x), \psi(x, y) = \psi(y, x), \phi(x, y) \neq \psi(x, y)$ and for arbitrary $\{x, y\} \in B, \phi(x, y) = x, \psi(x, y) = x$. An operation $m \in H$ is called arithmetical on $B$, if $\forall \{x, y\} \in B m(x, x, y) = m(y, x, x) = m(y, x, y) = y$.

The following theorem is a simple consequence of the main result of [21].
Suppose a set of operations \( \phi, \psi, m \) such that
- \( \phi, \psi \) is a tournament pair on \( B \),
- \( m \) is arithmetical on \( B \),

then \( \Gamma \) is MinHom-tractable.

The following theorem is a generalization of the previous one.

**Theorem 7.** Let \( \Gamma \) be a set of multi-sorted relations over a collection of finite sets \( A = \{ A_1, \ldots, A_n \} \) containing all unary multi-sorted relations. Assume that \( B_i \subseteq \{ \{ a, b \} | a, b \in A_i, a \neq b \} \). If MPol \( ( \Gamma ) \) contains a multi-sorted operations \( \phi, \psi, m \) such that
- \( \phi^{A_i}, \psi^{A_i} \) is a tournament pair on \( B_i \),
- \( m^{A_i} \) is arithmetical on \( B_i \),

then \( \Gamma^A \) is MinHom-tractable.

**Proof.** Denote \( G = \bigcup_{i=1}^{n} A_i \). It is easy to see that we can define operations \( \phi', \psi' : G^2 \to G \) and \( m' : G^3 \to G \) such that \( \phi'|_{A_i} = \phi^{A_i}, \psi'|_{A_i} = \psi^{A_i}, m'|_{A_i} = m^{A_i} \), and \( \phi', \psi' \) is a tournament pair on \( \bigcup_{i=1}^{n} B_i \cup \{ \{ a, b \} | a \in A_i, b \in A_j, i \neq j \} \). Then \( \phi', \psi', m' \in Pol ( \Gamma^A ) \) which by Theorem 6 means that \( \Gamma^A \) is MinHom-tractable.

### 5 Structure of \( R \)-tractable algebras

**Definition 18.** Suppose a set of operations \( H \) over \( D \) is conservative and \( O \subseteq \{ (x, y) \mid x, y \in D, x \neq y \} \). A pair of binary operations \( \phi, \psi \in H \) is called a weak tournament pair on \( O \), if

(1) \( \forall a, b \) such that \( (a, b), (b, a) \in O : \frac{\phi}{b} \frac{\psi}{b} \text{ or } \frac{\phi}{b} \frac{\psi}{b} \);

(2) \( \forall a, b \) such that \( (a, b) \in O, (b, a) \notin O : \frac{\phi}{b} \frac{\psi}{b} \text{ or } \frac{\phi}{b} \frac{\psi}{b} \); 

(3) \( \forall a, b \) such that \( (a, b) \notin O, (b, a) \notin O : \phi|_{\{a,b\}}(x,y) = x, \psi|_{\{a,b\}}(x,y) = y \).

For a binary operation \( f \in F \) define \( \text{Com}(f) = \{ \{ a, b \} : f|_{\{a,b\}} \text{ is commutative} \} \). Consider any binary operation \( f_{\text{max}} \in F \) which has maximal set \( \text{Com}(f_{\text{max}}) \), i.e. there is no \( f \in F \) such that \( \text{Com}(f_{\text{max}}) \subseteq \text{Com}(f) \). Since for any binary operations \( a, b \), \( \text{Com}(a(x,y)) \cup \text{Com}(b(x,y)) = \text{Com}(a(b(x,y), b(y,x))) \), we conclude that \( \text{Com}(f_{\text{max}}) = \bigcup_{f \in F} \text{Com}(f) \). Define \( \text{Com}^o(f_{\text{max}}) = \{ (a, b) \mid (a, b) \in \text{Com}(f_{\text{max}}) \} \).

**Theorem 8.** If \( F \) satisfies the necessary local conditions for \( R, UG(R) \) is complete and \( T_{B^R} = (M^o \cap E(R), P) \) is bipartite, then there is a pair \( \phi, \psi \in F \) which is a weak tournament pair on \( \text{Com}^o(f_{\text{max}}) \cap E(R) \).
Proof. Let $M_1, M_2$ denote a partitioning of vertices of the bipartite graph $T^R_B$. Then, for every $(a,b), (c,d) \in M_1$, there is a function $\phi \in F : \phi \uparrow \downarrow \in \phi$. Let us prove by induction that for every $(a_1,b_1), (a_2,b_2), \ldots, (a_n,b_n) \in M_1$, there is a function $\phi : \phi : a_1 \downarrow a_2 \downarrow \ldots \downarrow a_n \downarrow b_1 \downarrow b_2 \downarrow \ldots \downarrow b_n \downarrow$. The base of induction $n = 2$ is obvious. Let $(a_1,b_1), (a_2,b_2), \ldots, (a_{n+1}, b_{n+1}) \in M_1$ be given. By the induction hypothesis, there are $\phi_1, \phi_2, \phi_3 \in F : \phi_1, \phi_2, \phi_3 \downarrow a \downarrow a \downarrow \ldots \downarrow a \downarrow b \downarrow b \downarrow \ldots \downarrow b \downarrow \phi_3$. Then, it is easy to see that $\phi_1, \phi_2, \phi_3$ satisfy the first property of a weak tournament pair on $\Com^o(f_{\max}) \cap E(R)$. Our goal is to construct a pair of operations that satisfy other two properties.

Consider operations $\psi_2(x, y) = f_{\max}(\phi_2(x, y), \phi_2(y, x))$ and $\psi_3(x, y) = f_{\max}(\psi(x, y), \phi_3(x, y))$. Since for any $B \in \Com(\phi_2)$, $\phi_2|_B = \phi_2|_B$ and $\psi_2|_B = \psi_2|_B$, we obtain that $\phi_2 \uparrow \psi_3 \downarrow \in \phi_2 \uparrow \psi_3 \downarrow$ satisfy the first property of a weak tournament pair on $\Com^o(f_{\max}) \cap E(R)$, too. Moreover, $\phi_2 \uparrow \psi_3 \downarrow$ are commutative on $\Com(f_{\max})$. For any two elements $a, b \in A$ such that $(a, b) \in \Com^o(f_{\max}) \cap E(R), (b, a) \notin \Com^o(f_{\max}) \cap E(R), \phi_2|_{\{a, b\}}, \psi_3|_{\{a, b\}}$ are both commutative. If $(a, b)$ is a vertex in $T^R_B$, then one of $a \uparrow \psi_3 \downarrow$, and therefore it holds for one of $\phi_2 \uparrow \psi_3 \downarrow$, too.

If $(a, b)$ is not a vertex in $T^R_B$, i.e. $\{a, b\} \notin M$, then $\phi \uparrow \psi_3 \downarrow \in \phi \uparrow \psi_3 \downarrow$, because $F|_{\{a, b\}}$ contains only one commutative operation. In both cases the second property of a weak tournament pair on $\Com^o(f_{\max}) \cap E(R)$ is satisfied.

In case that $(a,b), (b,a) \notin \Com^o(f_{\max}) \cap E(R)$, we see that $\{a, b\} \notin \Com(f_{\max})$ (the case when $\{a, b\} \in \Com(f_{\max}), (a, b), (b, a) \notin E(R)$ is impossible due to completeness of $U G(R)$) which means that $\phi_2 \psi_3, \psi_3 \phi_2 \in \Com^o(f_{\max}) \cap E(R)$.

The analogous statement can be proved for $M_2$. So it follows from the proof that there are binary operations $\phi', \psi' \in F$, such that $\forall (x, y) \in M_1 : a \phi' \downarrow$ and $\forall (x, y) \in M_2 : b \psi' \downarrow$.

If $(a, b), (b, a) \in \Com^o(f_{\max}) \cap E(R)$, then the necessary local conditions for $C$ give that $\{a, b\} \in M$. Moreover, $(a, b), (b, a) \in M \cap E(R)$ are always in different partitions of $T^R_B$. Consequently, $\phi', \psi'$ satisfy the first property of a weak tournament pair on $\Com^o(f_{\max}) \cap E(R)$. Our goal is to construct a pair of operations that satisfy other two properties.

Consider operations $\phi''(x, y) = f_{\max}(\phi'(x, y), \phi'(y, x))$ and $\psi''(x, y) = f_{\max}(\psi'(x, y), \psi'(y, x))$. Since for any $B \in \Com(\phi')$, $\phi'(|_B = \phi'(|_B$ and $\psi'(|_B = \psi'(|_B$, we obtain that $\phi'' \psi'' \downarrow$ satisfy the first property of a weak tournament pair on $\Com^o(f_{\max}) \cap E(R)$, too. Moreover, $\phi'' \psi''$ are commutative on $\Com(f_{\max})$. For any two elements $a, b \in A$ such that $(a, b) \in \Com^o(f_{\max}) \cap E(R), (b, a) \notin \Com^o(f_{\max}) \cap E(R), \phi''|_{\{a, b\}}, \psi''|_{\{a, b\}}$ are both commutative. If $(a, b)$ is a vertex in $T^R_B$, then one of $a \uparrow \psi'' \downarrow$, and therefore it holds for one of $\phi'' \psi''$, too.

If $(a, b)$ is not a vertex in $T^R_B$, i.e. $\{a, b\} \notin M$, then $\phi \uparrow \psi'' \downarrow \in \phi \uparrow \psi'' \downarrow$, because $F|_{\{a, b\}}$ contains only one commutative operation. In both cases the second property of a weak tournament pair on $\Com^o(f_{\max}) \cap E(R)$ is satisfied.

In case that $(a,b), (b,a) \notin \Com^o(f_{\max}) \cap E(R)$, we see that $\{a, b\} \notin \Com(f_{\max})$ (the case when $\{a, b\} \in \Com(f_{\max}), (a, b), (b, a) \notin E(R)$ is impossible due to completeness of $U G(R)$) which means that $\phi''(a,b), \psi''(a,b)$ are projections. Thus, a pair of operations $\phi(x, y) = \phi''(x, \phi''(y, x))$, $\psi(x, y) = \psi''(\psi''(y, x), y)$ satisfy all three properties of a weak tournament pair on $\Com^o(f_{\max}) \cap E(R)$.

**Theorem 9.** If $F$ satisfies the necessary local conditions for $R$ and $\Com(f_{\max}) \neq \emptyset$, then $F$ contains an arithmetical operation on $\Com(f_{\max})$.

Proof. Obviously, for every $B \in \Com(f_{\max})$, $F|_B$ cannot contain any commutative binary function. Therefore, every binary function in $F|_B$ is a projection.
For $B \in \text{Com}(f_{\text{max}})$, let $m^B$ be an arithmetical function on $B$; existence of this function follows from the necessary local conditions for $R$. Assume now that $\text{Com}(f_{\text{max}}) = \{(x_1, y_1), \ldots, (x_s, y_s)\}$. We prove by induction that for every $r \leq s$, $F$ contains a function $m_r : A^3 \to A$ that is arithmetical on $\{(x_i, y_i) | 1 \leq i \leq r\}$.

When $r = 1$, $m_1(x, y, z) = m^{(x, y_1)}(x, y, z)$ and the statement is obviously true. Suppose it is true for $r \leq k < s$ and that we have the function $m_k : A^3 \to A$. Let us prove the statement for $r = k + 1$. If $m_k$ is arithmetical on $\{(x_{k+1}, y_{k+1})\}$, then we define $m_{k+1} \triangleq m_k$ and the statement is proved. Otherwise, one of the following three statements is true:

$$\exists x, y \in \{x_{k+1}, y_{k+1}\} \ [m_k(x, x, y) \neq y],$$

$$\exists x, y \in \{x_{k+1}, y_{k+1}\} \ [m_k(y, x, x) \neq y],$$

$$\exists x, y \in \{x_{k+1}, y_{k+1}\} \ [m_k(y, y, x) \neq y].$$

Suppose the first case holds (the proof for other cases is analogous), i.e. $m_k(x_{k+1}, y_{k+1})(x, x, y)$ is the $x$-projection. It is easy to see that the function $m_{k+1}(x, y, z) = m_k(m^{(x_{k+1}, y_{k+1})}(x, y, z), m^{(x_{k+1}, y_{k+1})}(x, y, z), m_k(x, y, z))$ is arithmetical on $\{(x_i, y_i) | 1 \leq i \leq k + 1\}$.

Induction completed and it is clear that $m_k(x, y, z)$ satisfies the condition of theorem.

6 Proof of Theorem 5

Suppose we have some instance $(V, C, W)$ of $\text{MinHom}_R(\text{Inv}(F))$. Let $\text{Sol} = (V, C, W)$ denote the set of its solutions. By Theorem 8, there is a pair $\phi, \psi \in F$ which is a weak tournament pair on $\text{Com}(f_{\text{max}}) \cap E(R)$.

Until the end of this section, the pair $\phi, \psi$ will be fixed. For any $f, g : V \to A$, define $\phi(f, g), \psi(f, g) : V \to A$ such that $\phi(f, g)(v) = \phi(f(v), g(v))$, $\psi(f, g)(v) = \psi(f(v), g(v))$. Since any constraint relation of $C$ is preserved by $\phi, \psi$, we have $\phi(f, g), \psi(f, g) \in \text{Sol} = (V, C, W)$ for $f, g \in \text{Sol} = (V, C, W)$.

**Lemma 2.** For any measure $M(f) = \sum_{v \in V} \sum_{r \in R} w_{vr} r(f(v))$ we have

$$M(\phi(f, g)) + M(\psi(f, g)) \leq M(f) + M(g)$$

**Proof.** Let us prove that for any $a, b \in A$ and any $r \in R$: $r(a) + r(b) \geq r(\phi(a, b)) + r(\psi(a, b))$.

If $(a, b), (b, a) \in \text{Com}(f_{\text{max}}) \cap E(R)$ or $(a, b) \notin \text{Com}(f_{\text{max}})$, then $\{\phi(a, b), \psi(a, b)\} = \{a, b\}$, and we have an equality. If $(a, b) \in \text{Com}(f_{\text{max}}) \cap E(R)$, $(b, a) \notin E(R)$, then $\{\phi(a, b), \psi(a, b)\} = \{a, b\}$ or $\{\phi(a, b), \psi(a, b)\} = \{b\}$. In the first case we have an equality and in the second case we have:

$$\forall r \in R, r(a) + r(b) \geq 2r(b) \iff \forall r \in C r(b) \leq r(a) \iff (b, a) \notin E(R)$$
From this inequality, we conclude that for any measure $M(f) = \sum_{v \in \mathcal{V}} \sum_{r \in \mathcal{R}} w_{vr} \phi(f(v))$ we have $M(f) + M(g) \geq M(\phi(f, g)) + M(\psi(f, g))$

**Lemma 3.** Suppose $A_v = \{f(v) | f \in Sol(V, C, W)\}$ and there is $a, b \in A_v$ such that $\downarrow b, \downarrow c$. Then we have

$$\min_{f \in \text{Sol}(V, C, W), f(v) = b} M(f) \leq \min_{f \in \text{Sol}(V, C, W), f(v) = a} M(f)$$

**Proof.** Suppose that $f^* \in \text{Sol}(V, C, W), f^*(v) = a$ and $f^{**} \in \text{Sol}(V, C, W), f^{**}(v) = b$ such that

$$M(f^*) = \min_{f \in \text{Sol}(V, C, W), f(v) = a} M(f), M(f^{**}) = \min_{f \in \text{Sol}(V, C, W), f(v) = b} M(f)$$

Since $\phi(f^*, f^{**})(v) = b, \psi(f^*, f^{**})(v) = b$, then by Lemma 2

$$2M(f^{**}) \leq M(\phi(f^*, f^{**})) + M(\psi(f^*, f^{**})) \leq M(f^*) + M(f^{**}) \quad \Box$$

Now we are ready to describe polynomial-time algorithm for $\text{MinHom}(Inv(F))$, which conceptually follows the proof of Theorem 8.3 in [3]. First we compute sets $A_v = \{f(v) | f \in \text{Sol}(V, C, W)\}$ for every $v \in \mathcal{V}$. This could be done polynomial-time, because, by result of [2], $\text{CSP}(Inv(F))$ is polynomial-time solvable by 3-consistency algorithm: adding a constraint $(v, \{a\})$ to $\mathcal{C}$ we can find whether $a \in A_v$ or not. After this operation, we iteratively delete all elements $a$ from every $A_v$ such that there is $b \in A_v : \downarrow a, \downarrow b$. This operation will not increase a minimum of optimized functional due to Lemma 3. Afterwards, we update every constraint pair $((v_1, \ldots, v_m), \rho) \in \mathcal{C}$ by defining $\rho = \rho \cap A_{v_1} \times \cdots \times A_{v_m}$.

Consider a collection of sets $\mathcal{A} = \{A_v | v \in \mathcal{V}\}$. Obviously, $|\mathcal{A}| < 2^{|A|}$. For a pair $((v_1, \ldots, v_m), \rho) \in \mathcal{C}$, a predicate $\rho$ can be considered as multi-sorted predicate over $\mathcal{A}$ with signature $(A_{v_1}, \ldots, A_{v_m})$. Moreover, for any $p \in F$ this multi-sorted predicate preserves a polymorphism $\{p^B | B \in \mathcal{A}\}$ where $p^B = p|B$.

By Theorem 9, $F$ contains an operation $m$ which is arithmetical on $\text{Com}(f_{\text{max}})$. Consider multi-sorted polymorphisms $\{\phi^B | B \in \mathcal{A}\}, \{\psi^B | B \in \mathcal{A}\}, \{m^B | B \in \mathcal{A}\}$. It is easy to see that these polymorphisms satisfy conditions of Theorem 7 and the set of multi-sorted predicates $\{\rho | (s, \rho) \in \mathcal{C}\}$ is tractable. Therefore, we can polynomially solve our problem.

### 7 Directions for further work

As it was said in the introduction, $\text{MinHom}_R$ fits the framework of Valued CSP ($\text{VCSP}$) [3]. The following subproblems of $\text{VCSP}$ generalize $\text{MinHom}$ and are not investigated yet.
1. What happens with $\text{MinHom}_R$ when $\text{UG}(R)$ is not complete. We believe that Theorem 5 is true for this case too.

2. Suppose we have a finite valued constraint language $\Gamma$, i.e. a set of valued predicates over some finite domain set. If $\Gamma$ contains all unary valued predicates, we call $\text{VCSP}(\Gamma)$ a conservative $\text{VCSP}$. This name is motivated by the fact that in this case the multimorphisms (which is a generalization of polymorphisms for valued constraint languages [3]) of $\Gamma$ must consist of conservative functions. Since there is a well-known dichotomy for conservative CSPs [2], we suspect that there is a dichotomy for conservative $\text{VCSPs}$.

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