Towards Sample-Optimal Compressive Phase Retrieval with Sparse and Generative Priors

Zhaoqiang Liu  
National University of Singapore  
dcslizha@nus.edu.sg

Subhroshekhar Ghosh  
National University of Singapore  
subhrowork@gmail.com

Jonathan Scarlett  
National University of Singapore  
scarlett@comp.nus.edu.sg

Abstract

Compressive phase retrieval is a popular variant of the standard compressive sensing problem in which the measurements only contain magnitude information. In this paper, motivated by recent advances in deep generative models, we provide recovery guarantees with near-optimal sample complexity for phase retrieval with generative priors. We first show that when using i.i.d. Gaussian measurements and an $L$-Lipschitz continuous generative model with bounded $k$-dimensional inputs, roughly $O(k \log L)$ samples suffice to guarantee that any signal minimizing an amplitude-based empirical loss function is close to the true signal. Attaining this sample complexity with a practical algorithm remains a difficult challenge, and finding a good initialization for gradient-based methods has been observed to pose a major bottleneck. To partially address this, we further show that roughly $O(k \log L)$ samples ensure sufficient closeness between the underlying signal and any globally optimal solution to an optimization problem designed for spectral initialization (though finding such a solution may still be challenging). We also adapt this result to sparse phase retrieval, and show that $O(s \log n)$ samples are sufficient for a similar guarantee when the underlying signal is $s$-sparse and $n$-dimensional, matching an information-theoretic lower bound. While these guarantees do not directly correspond to a practical algorithm, we propose a practical spectral initialization method motivated by our findings, and experimentally observe performance gains over various existing spectral initialization methods for sparse phase retrieval.

1 Introduction

In this paper, we consider the (real-valued) phase retrieval problem, which aims to recover a signal $x \in \mathbb{R}^n$ from noisy magnitude-only measurements:

$$y_i = |\langle a_i, x \rangle| + \eta_i, \quad i = 1, 2, \ldots, m,$$

where $a_i \in \mathbb{R}^n$ is the $i$-th sensing vector, and $\eta_i$ represents additive noise. This problem arises naturally in areas such as diffraction imaging, X-ray crystallography, microscopy, optics and astronomy [8], where it is often difficult or even impossible to observe the linear measurements directly, and one can only record the magnitudes or intensities (squared magnitudes).

In many real applications, to reduce the required number of measurements, it is of interest to exploit structure in the signal being estimated. In particular, for applications related to signal processing and imaging, it is well-known that the underlying signal typically admits a sparse or approximately sparse
representation in some known basis [40]. Motivated by this, and considering the popularity of the standard compressive sensing (CS) problem [16], the sparse phase retrieval problem has attracted significant research interest.

Moreover, inspired by the successful applications of deep generative models in many fields [15], recently, a new perspective of CS has emerged, for which the sparsity assumption is replaced by a generative model assumption. That is, instead of assuming sparsity, the signal is assumed to be close to the range of a generative model [5]. In addition to the theoretical developments in [5], the authors also provide impressive numerical results showing that for some imaging applications, using generative priors can significantly reduce the required number of measurements (e.g., by a factor of 5 to 10) for recovering the signal up to a given accuracy. Follow-up works of [5] include [58, 22, 45, 67, 12, 29], just to name a few.

In this paper, we focus on providing recovery guarantees with near-optimal sample complexity bounds (e.g., optimal up to constant factors) for the compressive phase retrieval problem, considering both sparse and generative priors.

1.1 Related Work

In this subsection, we provide a summary of some relevant works, which can roughly be divided into (i) the general phase retrieval problem with no structural assumptions on the signal, (ii) sparse phase retrieval, and (iii) phase retrieval with generative models.

**General phase retrieval:** A wide range of approaches have been designed to solve the phase retrieval problem. Error-reduction approaches [18, 14] work well in practice, but they lack provable guarantees. Convex methods such as PhaseLift [9] and PhaseCut [62] lift the phase retrieval problem to a higher dimension, and typically suffer from high computation cost. Convex relaxations that operate in the natural domain of the signal are proposed in [19, 2]. However, these convex relaxation based methods are not empirically competitive against widely-used non-convex methods.

Following the seminal work of Netrapalli et al. [42], several works have studied theoretically-guaranteed non-convex optimization algorithms for phase retrieval [8, 10, 63, 72]. These algorithms start from a spectral initialization method, and then use an iterative algorithm (e.g., alternating minimization or gradient descent) to further decrease the approximation error. For the general phase retrieval problem, the optimal sample complexity of $O(n)$ can be achieved by certain non-convex methods [10, 63, 72]. Interestingly, optimal (up to a logarithmic factor) sample complexity guarantees have also been provided in the case of a random initialization [55, 11]. It is worth mentioning that some variants of the classic Wirtinger Flow (WF) algorithm [8, 10] and trust-region methods [55] consider squared measurements:

$$y_i = |\langle a_i, x \rangle|^2 + \epsilon_i, \quad i = 1, 2, \ldots, m, \quad (2)$$

and minimize the intensity based empirical loss function:

$$f_f(w) := \frac{1}{2} \|y - |Aw|^2\|^2_2, \quad (3)$$

where $A \in \mathbb{R}^{m \times n}$ is the sensing matrix with its $i$-th row being $a_i^T$, and $|\cdot|$ is applied element-wise.

Nonetheless, numerical results suggest that algorithms minimizing the amplitude based empirical risk (with the measurement model (1))

$$f(w) := \|y - |Aw|\|^2_2 \quad (4)$$

are usually more efficient in computation [63, 72, 54]. Based on this, we focus on the measurement model (1) and the associated loss function (4).

**Sparse phase retrieval:** A variety of algorithms have been devised for sparse phase retrieval [44, 53, 50]. In particular, there exist various non-convex optimization based algorithms, including thresholded/projected WF [7, 54], sparse truncated amplitude flow [64], compressive phase retrieval with alternating minimization [28], and sparse phase retrieval by hard iterative pursuit [6]. All of these approaches are analyzed under the assumption of using a sensing matrix with i.i.d. Gaussian entries. In addition, similar to non-convex methods for the general phase retrieval problem, all of these approaches start with a spectral initialization step, and then use an iterative algorithm to refine the initial estimate.
We highlight two particularly relevant works in more detail. To guarantee favorable global geometry which are widely adopted for phase retrieval.

Algorithms proposed in both works \[20, 52\] minimize the objective function directly over the latent variable \(z\) with assumptions of noiseless measurements and the existence of a good initial vector. Recovery with \(n\) measurements model (1), and the loss function considered in \[66\] is different from both (3) and (4), error \[47\]. Moreover, the assumption of differentiability makes it is not directly applicable to the link function comes at a price; namely, it is an open problem to handle signals with representation setting of using differentiable but unknown link functions. The weaker assumption of an unknown guarantees for phase retrieval with generative models can also be found in \[66\], within a more general ReLU neural network priors and untrained neural network priors have been provided in \[25, 27\], achieving the optimal sample complexity \(O(s \log n)\). However, designing a computationally efficient algorithm achieving the optimal sample complexity for \(s\)-sparse vectors remains an important open problem.

**Phase retrieval with generative models**: Phase retrieval using generative models has been studied in \[20, 25, 27, 52, 66, 1\]. Extensive numerical experiments for phase retrieval with generative models, as well as both Gaussian and coded diffraction pattern measurements, have been presented in \[52\]. Algorithms proposed in both works \[20, 52\] minimize the objective function directly over the latent variable \(z \in \mathbb{R}^k\) using gradient descent, where \(k\) is the latent dimension. The corresponding objective function is highly non-convex, and performing gradient descent directly over \(z\) limits the explorable solution space, which may lead to getting stuck in local minima.

We highlight two particularly relevant works in more detail. To guarantee favorable global geometry for gradient methods, the authors of \[20\] assume a ReLU neural network generative model with i.i.d. zero-mean Gaussian weights and no offsets. In addition, the neural network needs to be sufficiently expansive such that \(n_i \geq \Omega(n_{i-1} \log n_{i-1})\), where \(n_i\) represents the number of neurons in the \(i\)-th layer. Under such conditions, a sample complexity \(O(k d^2 \log n)\) is obtained, where \(d\) is the number of layers. In another related work studying an approximate message passing algorithm, the authors of \[1\] maintain i.i.d. Gaussian weights but slightly relax to general activation functions (not only ReLU) and \(n_i \geq \Omega(n_{i-1})\). Under these conditions, the high-dimensional regime is studied \((n, m, k \to \infty\) with \(m/n\) kept fixed), and an asymptotic analysis is given (not yet established as fully rigorous). We also note that \[1\] focuses on the case that \(z \sim P_{z}\) for some separable \(P_{z}\). Both \[1\] and \[20\] focus on the noiseless case, though \[1\] states that the extension to noisy phase retrieval is possible.

Recovery guarantees for non-convex optimization algorithms for phase retrieval with pre-trained ReLU neural network priors and untrained neural network priors have been provided in \[25, 27\], with assumptions of noiseless measurements and the existence of a good initial vector. Recovery guarantees for phase retrieval with generative models can also be found in \[66\], within a more general setting of using differentiable but unknown link functions. The weaker assumption of an unknown link function comes at a price; namely, it is an open problem to handle signals with representation error \[47\]. Moreover, the assumption of differentiability makes it is not directly applicable to the measurement model \[1\], and the loss function considered in \[66\] is different from both \[3\] and \[4\], which are widely adopted for phase retrieval.

### 1.2 Contributions

- We show that for compressive phase retrieval with i.i.d. Gaussian measurements, when the signal is close to the range of an \(L\)-Lipschitz continuous generative model with bounded \(k\)-dimensional inputs, roughly \(O(k \log L)\) samples suffice for ensuring the closeness between the signal and any vector minimizing the amplitude-based empirical loss function \(4\).

- To address the sample complexity barrier posed by spectral initialization, we propose a relevant optimization problem, and show that roughly \(O(k \log L)\) samples are sufficient to guarantee that any *globally optimal* solution of the optimization problem is close to the true signal. This suggests the plausibility of practical spectral initialization algorithms that are able to find such a global optimum, though we do not attempt to provide one that can provably do so.

- We adapt our analysis to the case of sparse phase retrieval, and show that roughly \(O(s \log n)\) samples suffice for a similar recovery guarantee, where \(s\) is the sparsity level and \(n\) is the ambient dimension. This matches the information-theoretic lower bound up to constant factors.
We use upper and lower case boldface letters to denote matrices and vectors respectively. We write $[N] = \{1, 2, \ldots, N\}$ for a positive integer $N$. We define the $\ell_2$-ball $B_2^N(r) := \{z \in \mathbb{R}^k : \|z\|_2 \leq r\}$, and the unit sphere $S^{n-1} := \{x \in \mathbb{R}^n : \|x\|_2 = 1\}$. We use $G$ to denote an $L$-Lipschitz continuous generative model from $B_2^N(r)$ to $\mathbb{R}^n$. For a set $B \subseteq B_2^N(r)$, we write $G(B) = \{G(z) : z \in B\}$. The sensing matrix $A \in \mathbb{R}^{m \times n}$ is assumed to have i.i.d. standard normal entries, i.e., $a_{ij} \sim \mathcal{N}(0, 1)$. For any $s \in \mathbb{R}^n$, we use $\bar{s} = \frac{s}{\|s\|_2}$ to denote the corresponding normalized vector. The support (set) of a vector is the index set of its non-zero entries. For any $X \in \mathbb{R}^{m \times n}$ and any index set $I \subseteq [m]$, $X_I$ denotes the $|I| \times n$ sub-matrix of $X$ that only keeps the rows indexed by $I$. We use $\|X\|_{2 \rightarrow 2}$ to denote the spectral norm of a matrix $X$. For any positive integer $N$, we use $I_N$ to denote the identity matrix in $\mathbb{R}^{N \times N}$. We use the generic notations $C$ and $C'$ to denote large positive constants, and we use $c$ and $c'$ to denote small positive constants; their values may differ from line to line.

1.3 Notation

We use upper and lower case boldface letters to denote matrices and vectors respectively. We write $[N] = \{1, 2, \ldots, N\}$ for a positive integer $N$. We define the $\ell_2$-ball $B_2^N(r) := \{z \in \mathbb{R}^k : \|z\|_2 \leq r\}$, and the unit sphere $S^{n-1} := \{x \in \mathbb{R}^n : \|x\|_2 = 1\}$. We use $G$ to denote an $L$-Lipschitz continuous generative model from $B_2^N(r)$ to $\mathbb{R}^n$. For a set $B \subseteq B_2^N(r)$, we write $G(B) = \{G(z) : z \in B\}$. The sensing matrix $A \in \mathbb{R}^{m \times n}$ is assumed to have i.i.d. standard normal entries, i.e., $a_{ij} \sim \mathcal{N}(0, 1)$. For any $s \in \mathbb{R}^n$, we use $\bar{s} = \frac{s}{\|s\|_2}$ to denote the corresponding normalized vector. The support (set) of a vector is the index set of its non-zero entries. For any $X \in \mathbb{R}^{m \times n}$ and any index set $I \subseteq [m]$, $X_I$ denotes the $|I| \times n$ sub-matrix of $X$ that only keeps the rows indexed by $I$. We use $\|X\|_{2 \rightarrow 2}$ to denote the spectral norm of a matrix $X$. For any positive integer $N$, we use $I_N$ to denote the identity matrix in $\mathbb{R}^{N \times N}$. We use the generic notations $C$ and $C'$ to denote large positive constants, and we use $c$ and $c'$ to denote small positive constants; their values may differ from line to line.

2 Amplitude-Based Loss Minimization for Generative Priors

In this section, we provide recovery guarantees with respect to optimal solutions of the amplitude based loss function $f_2$ for phase retrieval with generative models. We let the generative model $G : B_2^N(r) \to \mathbb{R}^n$ be any $L$-Lipschitz continuous function, and we suppose that the underlying signal $x$ is close to (but does not need to lie in) the range of $G$. We are interested in the case $k \ll n$ (i.e., relatively small latent dimension). The sensing matrix $A$ is assumed to have i.i.d. standard normal entries. Let $q \in \text{Range}(G)$ minimize (4) to within additive $\tau > 0$ of the optimum, i.e.,

$$\|y - |Aq||_2^2 \leq \min_{w \in \text{Range}(G)} \|y - |Aw||_2^2 + \tau,$$

and define $p \in \text{Range}(G)$ to be the point that is closest to the signal $x$, i.e.,

$$p = \arg \min_{w \in \text{Range}(G)} \|x - w\|_2.$$

With the above settings in place, we present the following theorem, proved in Appendix [3]

**Theorem 1.** Consider the observation model (1) with i.i.d. Gaussian measurements, and $p, q$ satisfying (5) with parameter $\tau$. Then, for any $\delta > 0$, we have that if $m = \Omega(k \log \frac{L}{\epsilon})$ with probability $1 - e^{-\Omega(m)}$, it holds that

$$\min \{\|q - x\|_2, \|q + x\|_2\} \leq O \left(\|p - x\|_2 + \frac{\|y\|_2 + \sqrt{\tau}}{\sqrt{m}} + \delta\right).$$

Since we assume that $x$ is close to the range of $G$, the representation error $\|p - x\|_2$ is small (compared to $\|x\|_2$). The optimization error $\tau$ is also ideally small. In addition, it is common to assume that $\frac{\|y\|_2}{\sqrt{m}} \leq \|x\|_2$ (e.g., see (2)), where $c > 0$ is also small. A typical $d$-layer neural network generative model has $\text{poly}(n)$-bounded weights in each layer, and thus its Lipschitz constant is $L = n^{O(d)}$. Then, the values of $r$ and $\delta$ can be as extreme as $r = n^{O(d)}$ and $\delta = \frac{1}{2^{r+1}}$ without affecting the final scaling. In this case, Theorem 1 reveals that roughly $O(k \log L)$ samples suffice to guarantee that any optimal solution of the amplitude-based loss function (4) is close to the true signal (up to an unavoidable possible global sign flip). By the analysis of information-theoretic lower bounds in [31][37], this scaling cannot be improved without further assumptions.

---

1. Here and subsequently, all $m = \Omega(\cdot)$ statements are assumed to have a large enough implied constant.
3 Spectral Initializations with Generative Priors

In Theorem 1, we provided recovery guarantees with respect to optimal solutions of the amplitude based loss function. However, it is not clear whether such a guarantee is achievable by practical algorithms.

A notable recent work is [25], showing that for a depth-$d$, width-$w$ ReLU neural network, roughly $O(kd \log w)$ samples suffice for their alternating phase projected gradient descent (APPGD) approach to obtain an accurate estimate of the signal, as long as there exists a proper initial vector $x^{(0)}$, in the sense that $\min \{ \| x - x^{(0)} \|_2, \| x + x^{(0)} \|_2 \} \leq c \| x \|_2$ for some small positive constant $c$. However, as we know from sparse phase retrieval, designing a practical algorithm to find a proper initial vector with near-optimal sample complexity can be more difficult than designing the subsequent iterative algorithm that refines the initial guess, and this still remains as an open problem. It is also important to note that [25] assumes accurate projections onto the range of the generative model, but this is not guaranteed in practice, due to the use of approximations [51, 48].

Existing initialization methods for sparse phase retrieval typically first estimate the support of the sparse signal, and then perform a power method on a sub-matrix corresponding to the estimated support set to calculate an initial vector. Without extra assumptions beyond sparsity, these methods have a suboptimal sample complexity of $O(s^2 \log n)$. For phase retrieval with generative models, the situation is even worse – there is no theoretically-guaranteed initialization method at all. To address this gap, we provide recovery guarantees for a spectral initialization method of phase retrieval with generative models. We emphasize that our guarantees only concern the global optima of a suitably-defined optimization problem; see [14] below. Similarly to [25], the caveat remains that practical solutions may fail to find such optima (e.g., due to inexact projections).

We assume that $x$ is in the range of $G : B_2^k(r) \to \mathbb{R}^n$ (again assumed to be $L$-Lipschitz continuous), and we assume that the noise terms $\eta_1, \eta_2, \ldots, \eta_m$ are bounded as follows:

$$\frac{\| \eta \|_2}{\sqrt{m}} \leq c_0 \| x \|_2,$$

where $c_0 > 0$ is a sufficiently small constant. This assumption states that the signal to noise ratio (SNR) is sufficiently large, and has also been made in relevant existing works (see, e.g., [72, Theorem 3]). Moreover, we assume that

$$\| \eta \|_\infty \leq c_1 \| x \|_2,$$

where $c_1 > 0$ is also a sufficiently small constant. This assumption is satisfied, for example, when (assuming (9)),

$$\| \eta \|_\infty = O \left( \frac{\| \eta \|_2}{\sqrt{m}} \right),$$

i.e., when none of the noise entries are unusually large. Again, similar assumptions have been made in existing works such as [10, Theorem 2], and [71, Theorem 2].

Let $\lambda$ be defined as

$$\lambda := \sqrt{\frac{\pi}{2}} \frac{1}{m} \sum_{i=1}^{m} y_i.$$ (11)

Using sub-Gaussian concentration [60, Proposition 5.10], $\lambda$ is close to $\| x \|_2$ with high probability given enough samples. In the following, we focus on estimating the normalized signal vector $\hat{x} := \frac{x}{\| x \|_2}$. For this purpose, similar to the idea in [35], we consider a normalized generative model $G : D \to S^{n-1}$, where $D := \{ z \in B_2^k(r) : \| G(z) \|_2 > R_{\min} \}$ for some $R_{\min} > 0$ \footnote{This matches the $O(k \log L)$ scaling upon substituting $L = O(w^d)$ for ReLU networks [5].} $S^{n-1}$ denotes the unit sphere in $\mathbb{R}^n$, and $G(z) = \frac{G(z)}{\| G(z) \|_2}$. Then, $\hat{G}$ is $L$-Lipschitz continuous with $L = \frac{L}{R_{\min}}$. 

\footnote{When $\eta$ is Gaussian with $\eta \sim N(0, \sigma^2 I_m)$, (10) is not guaranteed, but any “truncated Gaussian” does satisfy (10). In addition, to circumvent this issue under Gaussian noise, we may instead assume that $\sigma \sqrt{\log m} \leq c_0 \| x \|_2$ in (8) (and remove the assumption in (9)). We also note that in [10, 71], the authors consider an intensity based measurement model, and accordingly, the corresponding assumption is for $\frac{\| \eta \|_2}{\| x \|_2}$ instead of for $\frac{\| \eta \|_\infty}{\| x \|_2}$.}

\footnote{As discussed in [35], the dependence on $R_{\min}$ in the sample complexity is very mild. Under the typical scaling $L = n^{O(d)}$ for a $d$-layer neural network, the scaling laws remain unchanged even with $R_{\min} = \frac{1}{n^{1/d}}$.}
In the following, we adopt an idea from [72] of considering an empirical matrix with a truncation operation. We note that [72] only considered doing so for general phase retrieval; we are not aware of any work doing so for constrained settings (e.g., with sparse or generative priors), and we found that handling such settings required substantial additional effort. In more detail, we consider the following matrix defined with a truncation operation (see Remark 2 below for a discussion comparing to the non-truncated counterpart):

\[ V := \frac{1}{m} \sum_{i=1}^{m} y_i a_i a_i^T 1_{\{l < y_i < u\}}, \]

(12)

where \( l, u \) are positive constants satisfying \( 1 < l < u \), and \( \lambda \) is given in (11). Note that for the noiseless case with \( y_i = |\langle a_i, x \rangle| \), in accordance with the above-mentioned closeness between \( \lambda \) and \( \|x\|_2 \), it is useful to consider the following expectation:

\[ J := \mathbb{E} \left[ \frac{1}{m} \sum_{i=1}^{m} y_i a_i a_i^T 1_{\{l \leq y_i < u\}} \right] = \|x\|_2 (\gamma_0 I_n + \beta_0 \tilde{x}\tilde{x}^T), \]

(13)

where for \( q \sim \mathcal{N}(0, 1) \), we define \( \gamma_0 := \mathbb{E}[|g| 1_{\{|l| < |u|\}}] \) and \( \beta_0 := \mathbb{E}[|g|^3 1_{\{|l| < |u|\}}] - \gamma_0 \) (cf. Lemma 3 in Appendix C.2). Based on the observation that \( \tilde{x} \) is a leading eigenvector of \( J \) and the assumption \( x \in \text{Range}(g) \), we consider using the following optimization problem to find an \( \hat{x} \) that approximates \( x \):

\[ \hat{x} := \arg \max_{w \in \tilde{G}(\tilde{D})} \; w^T V w. \]

(14)

In addition, in order to take the norm of \( x \) into account, we set the the initial vector \( x^{(0)} \) as

\[ x^{(0)} = \lambda \hat{x}. \]

(15)

We show that roughly \( O(k \log L) \) samples suffice for ensuring the closeness between \( x^{(0)} \) and \( x \) (up to a global signal flip). Formally, we have the following.

**Theorem 2.** Consider the model \( \mathcal{M} \) with i.i.d. Gaussian measurements and noise satisfying (9)-(12), and let \( \hat{x} \) be as defined in (14). Suppose that \( \hat{x} \in \tilde{G}(\tilde{D}) \) and \( u > 1 + c_1 \), where \( c_1 \geq \frac{\eta}{\|x\|_2} \) appears in \( c \).

Given a sufficiently small positive constant \( c \), when \( m = \Omega(k \log(\tilde{L} \tilde{n} r)) \), we have with probability \( 1 - e^{-\Omega(m)} \) that

\[ \min\{\|\hat{x} - \hat{x}\|_2, \|\hat{x} + \hat{x}\|_2\} < 0.9c. \]

(16)

In addition, we have with probability \( 1 - e^{-\Omega(m)} \) that \( x^{(0)} = \lambda \hat{x} \) satisfies

\[ \min\{\|x^{(0)}\|_2, \|x^{(0)} + x\|_2\} < c\|x\|_2. \]

(17)

The proof is given in Appendix C. Using well-established chaining arguments as those in [5, 35, 36], the sample complexity \( \Omega(k \log (\tilde{L} \tilde{n} r)) \) can be reduced to \( \Omega(k \log (\tilde{L} \tilde{r})) \). However, since for a typical \( d \)-layer neural network, we have the Lipschitz constant \( L = n^{O(d)} \) (and thus \( \tilde{L} \) and \( r \) can also be of order \( n^{O(d)} \) [5]), such a reduction is typically of minor importance.

## 4 Sparse Phase Retrieval

The proof of Theorem 2 relies on the fact that \( \hat{x} \) lies in \( \text{Range}(\tilde{G}) \subseteq S^{n-1} \), with the \( \delta \)-covering number being upper bounded by \( \left( \frac{L \delta}{\lambda} \right)^{k} \). Letting \( \Sigma^n_s \) be the set of all \( s \)-sparse unit vectors in \( \mathbb{R}^n \), we know that the \( \delta \)-covering number of \( \Sigma^n_s \) is upper bounded by \( \left( \frac{s}{\lambda} \right)^{k} \left( \frac{L \delta}{\lambda} \right)^{k} \leq \left( \frac{s}{\lambda} \right)^{k} \). Based on this observation, we can readily adapt the proof of Theorem 2 to obtain the following theorem.

---

5This is weaker than the assumption that \( \|x\|_2 > R_{\min} \) made in [35]. Moreover, even if \( \|x\|_2 \leq R_{\min} \), the condition \( \hat{x} \in \tilde{G}(\tilde{D}) \) may be satisfied (e.g., when \( G \) is a ReLU neural network generative model with no offsets, then the range of \( G \) becomes a cone, and if \( 0 < \|x\|_2 \leq R_{\min} \), we can choose a sufficiently large \( C \) such that \( C\|x\|_2 > R_{\min} \), with \( Ca \in \text{Range}(G) \).

6Note that the smaller \( c \) is taken to be here, the smaller the constant \( c_0 \) needs to be in the assumption (6).
We will see in our experiments that the method corresponding to the truncated empirical matrix $V$ outperforms the method corresponding to the non-truncated empirical matrix $\tilde{V}$, although the expectations of these two matrices have a very similar structure. Thus, we believe that truncation may potentially help in overcoming the computational challenges discussed above.

We note that [43, 56] consider solving sparse phase retrieval via practical SPCA methods, in a more general setting considering an unknown link function. They consider a weighted empirical covariance matrix with no truncation, and the derived sample complexity has a quadratic dependence on $s$. 

**Theorem 3.** Consider the model [1] with i.i.d. Gaussian measurements and noise satisfying (5–9), and suppose that $x \in \mathbb{R}^n$ is $s$-sparse. Let $\lambda$ be defined as in (11). Suppose that $u > l > 1 + c_1$, where $c_1 \geq \frac{\|\eta\|}{\|x\|_2}$ appears in [9], and $V$ is defined in (12). Let $\hat{x}$ be defined as

$$\hat{x} := \arg \max_{w \in \Sigma_2} w^T V w. \quad (18)$$

Then, when $m = \Omega(s \log n)$, we have with probability $1 - e^{-\Omega(m)}$ that

$$\min\{\|\hat{x} - x\|_2, \|\hat{x} + x\|_2\} < 0.9c,$$  

(19)

where $c > 0$ is a sufficiently small constant. In addition, we have with probability $1 - e^{-\Omega(m)}$ that $x^{(0)} = \lambda \hat{x}$ satisfies

$$\min\{\|x^{(0)} - x\|_2, \|x^{(0)} + x\|_2\} < c\|x\|_2.$$

(20)

**Implications regarding SPCA and support recovery:** We note that (18) is essentially a problem of sparse principal component analysis (SPCA) [7, 33]. Existing spectral initialization methods for sparse phase retrieval typically start from estimating the support of $x$, with an accurate estimate of the support requiring a suboptimal $O(s^2 \log n)$ sample complexity [7, 64, 28]. This serves as the major bottleneck of the sample complexity upper bound of the whole procedure (i.e., the initialization step and the subsequent iterative step), and it is unknown whether there is a practical and sample-optimal approach for estimating the support. In fact, to our knowledge, even the following question has not been answered previously:

**Key question:** Does there exist any spectral-based method (computationally efficient or otherwise) attaining an accurate initial vector with near-optimal sample complexity?

We provide a positive answer to this question, showing that $O(s \log n)$ samples suffice to ensure that any optimal solution of the SPCA problem (18) is close to the underlying (normalized) signal.

Motivated by the strong sample complexity guarantee provided in Theorem 3, we propose a spectral initialization method for sparse phase retrieval as described in Algorithm 1. In particular, after calculating $V$ from $A$ and $y$, we use algorithms for SPCA to solve (18). We observe via numerical experiments in Section 5 that our initialization method outperforms several popular spectral initialization methods used for sparse phase retrieval.

**Remark 1.** In various works on SPCA, it has been shown that there exists a statistical-computational gap for achieving consistency and optimal convergence rates uniformly over a parameter space (e.g., see [42, 65, 27]). However, the models therein are all different from ours (e.g., the spiked covariance model [17]), so it is unclear to what extent we should expect similar computational difficulties.

**Remark 2.** The truncation function $1_{\{(l < y_i < u)\}}$ in (12) plays an important role in our analysis. Specifically, it allows us to apply sub-exponential concentration bounds to certain quantities of interest, where the non-truncated version would involve heavier-tailed random variables (e.g., sub-Weibull of order $\alpha < 1$ [27, Theorem 3.1]). We believe that this distinction could be even more important when seeking uniform recovery guarantees, though we do not do so in this work. In Section 5 we present numerical results for applying SPCA to the non-truncated weighted empirical matrix

$$\tilde{V} = \frac{1}{m} \sum_{i=1}^m y_i a_i a_i^T. \quad (21)$$

We will see in our experiments that the method corresponding to the truncated empirical matrix $V$ outperforms the method corresponding to the non-truncated empirical matrix $\tilde{V}$, although the expectations of these two matrices have a very similar structure. Thus, we believe that truncation may potentially help in overcoming the computational challenges discussed above.
We compare our method (PRI-SPCA) with spectral initialization methods used in several popular algorithms for sparse phase retrieval: Thresholded Wirtinger flow (ThWF) \cite{7}, sparse truncated amplitude flow (SPARTA) \cite{64}, and compressive phase retrieval with alternating minimization (CoPRAM) \cite{23}. We also compare with the approach of solving an SPCA problem with respect to \( V \) (cf. \cite{21}), which is a non-truncated version of \( V \). The corresponding method is termed PRI-SPCA-NT, with ‘NT’ standing for no truncation. For PRI-SPCA, we set \( l = 1 \) and \( u = 5 \), as previously suggested in \cite{72}.

Various algorithms have been developed to tackle the SPCA problem, see, e.g., \cite{41, 30, 39}. In our experiments, we present the numerical results using the generalized Rayleigh quotient iteration (GRQI) method \cite{33}, but we observed that other popular SPCA methods such as the truncated power method (TPower) \cite{70} give similar averaged relative errors. For GRQI, we set the deflation parameter to be 0.2, and the total number of iterations to be 100. Since the sparsity level \( s \) is typically assumed to be known a priori for SPARTA and CoPRAM (ThWF assumes the knowledge of a parameter that plays the similar role as \( s \)), for GRQI, we set the parameter about the maximum number of non-zero indices as \( s \). Note that the time complexity of each iteration of GRQI is \( O(ns + s^3) \), which scales mildly compared to the subsequent \( O(s^2n\log n) \) time complexity of the iterative methods used in the three approaches (for ThWF, the corresponding time complexity is \( O(n^2\log n) \). See \cite[Table I]{28}). The number of power method steps used in all the methods is fixed to be 100.

The authors of \cite{33} suggest using the column of maximal norm as a starting point of GRQI. Because this policy has a time complexity of \( O(n^2) \), we instead choose the column corresponding to the largest diagonal entry as the starting point. Due to the structure of \( V \) (cf. \cite{13}), this modified strategy is very similar to the suggested policy, but with a smaller time complexity requirement. For a fair comparison, all power methods used in ThWF, SPARTA, and CoPRAM are similarly initialized. For the initialization method used in ThWF, we set the tuning parameter as \( \alpha = 0.1 \), as suggested in \cite{7}. Similarly, we set [\( I = \lceil \frac{1}{6}m \rceil \)] for the initialization method of SPARTA.

The support of the signal vector \( x \) is uniformly random, and the non-zero entries are generated from i.i.d. standard normal distributions. The noise vector \( \eta \) is generated from the distribution \( \mathcal{N}(0, \sigma^2\|x\|_2^2I_m) \), where \( \sigma > 0 \) represents the noise level. The initial relative error is defined as

\[
\min \{ \|x - x^{(0)}\|_2, \|x + x^{(0)}\|_2 \} / \|x\|_2,
\]

where \( x^{(0)} \) is an initial vector produced by each of the above mentioned initialization methods. The reported relative error is averaged over 50 random trials, and error bars indicate a single standard deviation over these trials. All numerical experiments are conducted using MATLAB R2014b on a machine with an Intel Core i5 CPU at 1.8 GHz and 8GB RAM.

\footnote{We use the MATLAB package shared by the authors of \cite{28} at \url{https://github.com/GauriJagatap/model-copram}. The package not only contains the codes for CoPRAM, but also for ThWF and SPARTA.}

\footnote{Using TPower also results in small relative error for our PRI-SPCA method, but in general, we found it to be less stable than GRQI in the absence of a careful design of its starting point.}
1. **Sample size effect:** In our first experiment, similarly to [64], we consider noiseless measurements, and we fix the signal dimension as \( n = 1000 \). We consider \( s \) equaling both 10 and 20. The number of measurements \( m \) takes values in \( \{100, 200, \ldots, 3000\} \). From Figure 1 we observe that our PRI-SPCA method almost always achieves the smallest initial relative error, thus outperforming all other methods, including the PRI-SPCA-NT method, which uses the non-truncated empirical matrix \( \tilde{V} \).

2. **Sparsity effect:** We consider noiseless measurements and fix \( n = 1000 \), and set \( m = 1000 \) or \( 2000 \), and vary \( s \) in \( \{5, 10, \ldots, 50\} \). We observe from Figure 2 that for this experiment, our PRI-SPCA method also gives the best empirical performance.

3. **Noise effect:** To demonstrate that PRI-SPCA also performs well in the noisy case, we vary \( \sigma \) in \( \{0.1, 0.2, \ldots, 1\} \). In addition, similarly to [64], we fix \( n = 1000 \), \( m = 3000 \), and consider \( s \) equaling both 10 and 20. For this experiment, similarly to that in [28], we do not compare with ThWF, as it uses quadratic measurements. From Figure 3 we observe that PRI-SPCA outperforms SPARTA and CoPRAM, and it is better than PRI-SPCA-NT at low noise levels.

In Appendix E we additionally compare with other initialization methods in terms of the effectiveness when combining with the subsequent iterative algorithm of CoPRAM. In particular, we compare the relative error and empirical success rate in the noiseless case, whereas for the noisy case, we compare the relative error when using approximately the same running time. We observe that for most cases, our PRI-SPCA method is superior to all other methods we compare with. This again suggests that PRI-SPCA can serve as a good initialization method in practice.

### 6 Conclusion and Future Work

We have provided a near-optimal recovery guarantee for amplitude-based loss minimization for phase retrieval with generative priors. In addition, motivated by the bottleneck of the spectral initialization for compressive phase retrieval, we have provided near-optimal recovery guarantees with respect to the optimal solutions of optimization problems designed for spectral initialization. An immediate future research direction is to design provably sample-optimal and computationally efficient spectral initialization methods that build on our theoretical results. In addition, extensions to complex-valued phase retrieval [32][46] would be of significant interest.

**Acknowledgment.** This work was supported by the Singapore National Research Foundation (NRF) under grant R-252-000-A74-281, and the Singapore Ministry of Education (MoE) under grants R-146-000-250-133 and R-146-000-312-114.
Figure 1: Average relative error vs. number of measurements $m$ with $s = 10$ (Left), $s = 20$ (Right).

Figure 2: Average relative error vs. sparsity level $s$ with $m = 1000$ (Left), $m = 2000$ (Right).

Figure 3: Average relative error vs. noise level $\sigma$ with $s = 10$ (Left), $s = 20$ (Right).
References

[1] B. Aubin, B. Loureiro, A. Baker et al., “Exact asymptotics for phase retrieval and compressed sensing with random generative priors,” in *Math. Sci. Mach. Learn. (MSML)*. PMLR, 2020, pp. 55–73.

[2] S. Bahmani and J. Romberg, “Phase retrieval meets statistical learning theory: A flexible convex relaxation,” in *Int. Conf. Artif. Intell. Stat. (AISTATS)*. PMLR, 2017, pp. 252–260.

[3] R. Baraniuk, M. Davenport, R. DeVore, and M. Wakin, “A simple proof of the restricted isometry property for random matrices,” *Constr. Approx.*, vol. 28, no. 3, pp. 253–263, 2008.

[4] Q. Berthet and P. Rigollet, “Complexity theoretic lower bounds for sparse principal component detection,” in *Conf. Learn. Theory (COLT)*. PMLR, 2013, pp. 1046–1066.

[5] A. Bora, A. Jalal, E. Price, and A. G. Dimakis, “Compressed sensing using generative models,” in *Int. Conf. Mach. Learn. (ICML)*, 2017, pp. 537–546.

[6] J.-F. Cai, J. Li, X. Lu, and J. You, “Sparse signal recovery from phaseless measurements via hard thresholding pursuit,” [https://arxiv.org/abs/2005.08777](https://arxiv.org/abs/2005.08777), 2020.

[7] T. T. Cai, X. Li, Z. Ma et al., “Optimal rates of convergence for noisy sparse phase retrieval via thresholded Wirtinger flow,” *Ann. Stat.*, vol. 44, no. 5, pp. 2221–2251, 2016.

[8] E. J. Candès, X. Li, and M. Soltanolkotabi, “Phase retrieval via Wirtinger flow: Theory and algorithms,” *IEEE Trans. Inf. Theory*, vol. 61, no. 4, pp. 1985–2007, 2015.

[9] E. J. Candès, T. Strohmer, and V. Voroninski, “PhaseLift: Exact and stable signal recovery from magnitude measurements via convex programming,” *Comm. Pure Appl. Math.*, vol. 66, no. 8, pp. 1241–1274, 2013.

[10] Y. Chen and E. J. Candès, “Solving random quadratic systems of equations is nearly as easy as solving linear systems,” *Comm. Pure Appl. Math.*, vol. 70, no. 5, pp. 822–883, 2017.

[11] Y. Chen, Y. Chi, J. Fan, and C. Ma, “Gradient descent with random initialization: Fast global convergence for nonconvex phase retrieval,” *Math. Program.*, vol. 176, no. 1, pp. 5–37, 2019.

[12] J. Cocola, P. Hand, and V. Voroninski, “Nonasymptotic guarantees for spiked matrix recovery with generative priors,” in *Conf. Neur. Inf. Proc. Sys. (NeurIPS)*, vol. 33, 2020.

[13] Y. C. Eldar and S. Mendelson, “Phase retrieval: Stability and recovery guarantees,” *Appl. Comput. Harmon. Anal.*, vol. 36, no. 3, pp. 473–494, 2014.

[14] J. R. Fienup, “Phase retrieval algorithms: A comparison,” *Appl. Opt.*, vol. 21, no. 15, pp. 2758–2769, 1982.

[15] D. Foster, *Generative Deep Learning : Teaching Machines to Paint, Write, Compose and Play*. O’Reilly Media, Inc. USA, 2019.

[16] S. Foucart and H. Rauhut, *A Mathematical Introduction to Compressive Sensing*. Springer New York, 2013.

[17] C. Gao, Z. Ma, H. H. Zhou et al., “Sparse CCA: Adaptive estimation and computational barriers,” *Ann. Stat.*, vol. 45, no. 5, pp. 2074–2101, 2017.

[18] R. W. Gerchberg, “A practical algorithm for the determination of phase from image and diffraction plane pictures,” *Optik*, vol. 35, pp. 237–246, 1972.

[19] T. Goldstein and C. Studer, “Phasemax: Convex phase retrieval via basis pursuit,” *IEEE Trans. Inf. Theory*, vol. 64, no. 4, pp. 2675–2689, 2018.

[20] P. Hand, O. Leong, and V. Voroninski, “Phase retrieval under a generative prior,” in *Conf. Neur. Inf. Proc. Sys. (NeurIPS)*, 2018, pp. 9154–9164.

[21] B. Hao, Y. Yadkori, Z. Wen, and G. Cheng, “Bootstrapping upper confidence bound,” *Conf. Neur. Inf. Proc. Sys. (NeurIPS)*, 2019.

[22] R. Heckel and P. Hand, “Deep decoder: Concise image representations from untrained non-convolutional networks,” in *Int. Conf. Learn. Repr. (ICLR)*, 2019.

[23] M. Huang and Z. Xu, “The estimation performance of nonlinear least squares for phase retrieval,” *IEEE Trans. Inf. Theory*, vol. 66, no. 12, pp. 7967–7977, 2020.

[24] M. Huang and Z. Xu, “Performance bound of the intensity-based model for noisy phase retrieval,” [https://arxiv.org/abs/2004.08764](https://arxiv.org/abs/2004.08764), 2020.
[25] R. Hyder, V. Shah, C. Hegde, and M. S. Asif, “Alternating phase projected gradient descent with generative priors for solving compressive phase retrieval,” in *Int. Conf. Acoust. Sp. Sig. Proc. (ICASSP)*. IEEE, 2019, pp. 7705–7709.

[26] M. Iwen, A. Viswanathan, and Y. Wang, “Robust sparse phase retrieval made easy,” *Appl. Comput. Harmon. Anal.*, vol. 1, no. 42, pp. 135–142, 2017.

[27] G. Jagatap and C. Hegde, “Algorithmic guarantees for inverse imaging with untrained network priors,” in *Conf. Neur. Inf. Proc. Sys. (NeurIPS)*, vol. 32, 2019.

[28] G. Jagatap and C. Hegde, “Sample-efficient algorithms for recovering structured signals from magnitude-only measurements,” *IEEE Trans. Inf. Theory*, vol. 65, no. 7, pp. 4434–4456, 2019.

[29] A. Jalal, L. Liu, A. G. Dimakis, and C. Caramanis, “Robust compressed sensing using generative models,” *Conf. Neur. Inf. Proc. Sys. (NeurIPS)*, 2020.

[30] M. Journée, Y. Nesterov, P. Richtárik, and R. Sepulchre, “Generalized power method for sparse principal component analysis.” *J. Mach. Learn. Res.*, vol. 11, no. 2, 2010.

[31] A. Kamath, S. Karmalkar, and E. Price, “On the power of compressed sensing with generative models.” in *Int. Conf. Mach. Learn. (ICML)*, 2020.

[32] F. Krahmer and D. Stöger, “Complex phase retrieval from subgaussian measurements,” *J. Fourier Anal. Appl.*, vol. 26, no. 6, pp. 1–27, 2020.

[33] V. Kuleshov, “Fast algorithms for sparse principal component analysis based on Rayleigh quotient iteration,” in *Int. Conf. Mach. Learn. (ICML)*. PMLR, 2013, pp. 1418–1425.

[34] G. Lecué and S. Mendelson, “Minimax rate of convergence and the performance of empirical risk minimization in phase recovery,” *Electron. J. Probab.*, vol. 20, pp. 1–29, 2015.

[35] Z. Liu, S. Gomes, A. Tiwari, and J. Scarlett, “Sample complexity bounds for 1-bit compressive sensing and binary stable embeddings with generative priors,” in *Int. Conf. Mach. Learn. (ICML)*, 2020.

[36] Z. Liu and J. Scarlett, “The generalized Lasso with nonlinear observations and generative priors,” in *Conf. Neur. Inf. Proc. Sys. (NeurIPS)*, vol. 33, 2020.

[37] Z. Liu and J. Scarlett, “Information-theoretic lower bounds for compressive sensing with generative models,” *IEEE J. Sel. Areas Inf. Theory*, vol. 1, no. 1, pp. 292–303, 2020.

[38] T. Ma and A. Wigderson, “Sum-of-squares lower bounds for sparse PCA,” *Conf. Neur. Inf. Proc. Sys. (NeurIPS)*, vol. 28, pp. 1612–1620, 2015.

[39] Z. Ma, “Sparse principal component analysis and iterative thresholding,” *Ann. Stat.*, vol. 41, no. 2, pp. 772–801, 2013.

[40] S. Mallat, *A wavelet tour of signal processing*. Elsevier, 1999.

[41] B. Moghaddam, Y. Weiss, and S. Avidan, “Spectral bounds for sparse PCA: Exact and greedy algorithms,” *Conf. Neur. Inf. Proc. Sys. (NeurIPS)*, vol. 18, p. 915, 2006.

[42] P. Netrapalli, P. Jain, and S. Sanghavi, “Phase retrieval using alternating minimization,” *IEEE Trans. Sig. Proc.*, vol. 63, no. 18, pp. 4814–4826, 2015.

[43] M. Neykov, Z. Wang, and H. Liu, “Agnostic estimation for misspecified phase retrieval models,” *Conf. Neur. Inf. Proc. Sys. (NeurIPS)*, 2016.

[44] H. Ohlsson, A. Yang, R. Dong, and S. Sastry, “CPRL--An extension of compressive sensing to the phase retrieval problem,” *Conf. Neur. Inf. Proc. Sys. (NeurIPS)*, vol. 25, pp. 1367–1375, 2012.

[45] G. Ongie, A. Jalal, C. A. Metzler, R. G. Baraniuk, A. G. Dimakis, and R. Willett, “Deep learning techniques for inverse problems in imaging,” *IEEE J. Sel. Areas Inf. Theory*, vol. 1, no. 1, pp. 39–56, 2020.

[46] T. Pang, Q. Li, Z. Wen, and Z. Shen, “Phase retrieval: A data-driven wavelet frame based approach,” *Appl. Comput. Harmon. Anal.*, vol. 49, no. 3, pp. 971–1000, 2020.

[47] Y. Plan and R. Vershynin, “The generalized Lasso with non-linear observations,” *IEEE Trans. Inf. Theory*, vol. 62, no. 3, pp. 1528–1537, 2016.

[48] A. Raj, Y. Li, and Y. Bresler, “GAN-based projector for faster recovery with convergence guarantees in linear inverse problems,” in *IEEE/CVF Int. Conf. Comp. Vision (ICCV)*, 2019.
[49] J. Scarlett and V. Cevher, “An introductory guide to Fano’s inequality with applications in statistical estimation,” https://arxiv.org/1901.00555, 2019.

[50] P. Schniter and S. Rangan, “Compressive phase retrieval via generalized approximate message passing,” IEEE Trans. Sig. Proc., vol. 63, no. 4, pp. 1043–1055, 2014.

[51] V. Shah and C. Hegde, “Solving linear inverse problems using GAN priors: An algorithm with provable guarantees,” in Int. Conf. Acoust. Sp. Sig. Proc. (ICASSP). IEEE, 2018, pp. 4609–4613.

[52] F. Shamshad and A. Ahmed, “Compressed sensing-based robust phase retrieval via deep generative priors,” IEEE Sens. J., vol. 21, no. 2, pp. 2286–2298, 2020.

[53] Y. Shechtman, A. Beck, and Y. C. Eldar, “GESPAR: Efficient phase retrieval of sparse signals,” IEEE Trans. Sig. Proc., vol. 62, no. 4, pp. 928–938, 2014.

[54] M. Soltanolkotabi, “Structured signal recovery from quadratic measurements: Breaking sample complexity barriers via nonconvex optimization,” IEEE Trans. Inf. Theory, vol. 65, no. 4, pp. 2374–2400, 2019.

[55] J. Sun, Q. Qu, and J. Wright, “A geometric analysis of phase retrieval,” Found. Comput. Math., vol. 18, no. 5, pp. 1131–1198, 2018.

[56] Y. S. Tan, “Sparse phase retrieval via sparse PCA despite model misspecification: A simplified and extended analysis,” https://arxiv.org/1712.04106, 2017.

[57] L. V. Truong and J. Scarlett, “Support recovery in the phase retrieval model: Information-theoretic fundamental limits,” IEEE Trans. Inf. Theory, vol. 66, no. 12, pp. 7887–7910, 2020.

[58] D. Van Veen, A. Jalal, M. Soltanolkotabi, E. Price, S. Vishwanath, and A. G. Dimakis, “Compressed sensing with deep image prior and learned regularization,” https://arxiv.org/1806.06438, 2018.

[59] S. S. Vempala, The random projection method. American Mathematical Soc., 2005, vol. 65.

[60] R. Vershynin, “Introduction to the non-asymptotic analysis of random matrices,” https://arxiv.org/abs/1011.3027, 2010.

[61] V. Vovkonski and Z. Xu, “A strong restricted isometry property, with an application to phaseless compressed sensing,” Appl. Comput. Harmon. Anal., vol. 40, no. 2, pp. 386–395, 2016.

[62] I. Waldspurger, A. d’Aspremont, and S. Mallat, “Phase recovery, maxcut and complex semidefinite programming,” Math. Program., vol. 149, no. 1, pp. 47–81, 2015.

[63] G. Wang, G. B. Giannakis, and Y. C. Eldar, “Solving systems of random quadratic equations via truncated amplitude flow,” IEEE Trans. Inf. Theory, vol. 64, no. 2, pp. 773–794, 2017.

[64] G. Wang, L. Zhang, G. B. Giannakis, M. Akçakaya, and J. Chen, “Sparse phase retrieval via truncated amplitude flow,” IEEE Trans. Sig. Proc., vol. 66, no. 2, pp. 479–491, 2017.

[65] T. Wang, Q. Berthet, R. J. Samworth et al., “Statistical and computational trade-offs in estimation of sparse principal components,” Ann. Stat., vol. 44, no. 5, pp. 1896–1930, 2016.

[66] X. Wei, Z. Yang, and Z. Wang, “On the statistical rate of nonlinear recovery in generative models with heavy-tailed data,” in Int. Conf. Mach. Learn. (ICML), 2019, pp. 6697–6706.

[67] J. Whang, Q. Lei, and A. G. Dimakis, “Compressed sensing with invertible generative models and dependent noise,” https://arxiv.org/2003.08089, 2020.

[68] F. Wu and P. Rebeschini, “Hadamard Wirtinger flow for sparse phase retrieval,” in Int. Conf. Artif. Intell. Stat. (AISTATS). PMLR, 2021, pp. 982–990.

[69] Y. Xia and Z. Xu, “The recovery of complex sparse signals from few phaseless measurements,” Appl. Comput. Harmon. Anal., vol. 50, pp. 1–15, 2021.

[70] X.-T. Yuan and T. Zhang, “Truncated power method for sparse eigenvalue problems,” J. Mach. Learn. Res., vol. 14, no. 4, 2013.

[71] H. Zhang, Y. Chi, and Y. Liang, “Provable non-convex phase retrieval with outliers: Median truncated Wirtinger flow,” in Int. Conf. Mach. Learn. (ICML). PMLR, 2016, pp. 1022–1031.

[72] H. Zhang, Y. Zhou, Y. Liang, and Y. Chi, “A nonconvex approach for phase retrieval: Reshaped Wirtinger flow and incremental algorithms,” J. Mach. Learn. Res., vol. 18, 2017.

[73] H. Zou, T. Hastie, and R. Tibshirani, “Sparse principal component analysis,” J. Comput. Graph. Stat., vol. 15, no. 2, pp. 265–286, 2006.
A General Auxiliary Results

In this section, we provide some useful auxiliary results that will be used throughout our analysis. First, we have the following basic concentration inequality.

Lemma 1. ([59, Lemma 1.3]) Suppose that $A \in \mathbb{R}^{m \times n}$ has i.i.d. standard normal entries. For fixed $x \in \mathbb{R}^n$, we have for any $\epsilon \in (0, 1)$ that

$$P\left( (1 - \epsilon)\|x\|_2^2 \leq \left\| \frac{1}{\sqrt{m}} Ax \right\|_2^2 \leq (1 + \epsilon)\|x\|_2^2 \right) \geq 1 - 2e^{-\epsilon^2(1-\epsilon)^m/4}. \quad (23)$$

Next, we state the following standard definition.

Definition 1. A random variable $X$ is said to be sub-Gaussian if there exists a positive constant $C$ such that $(\mathbb{E}[|X|^p])^{1/p} \leq C\sqrt{p}$ for all $p \geq 1$. The sub-Gaussian norm of a sub-Gaussian random variable $X$ is defined as $\|X\|_{\psi_2} = \sup_{p \geq 1} p^{-1/2} (\mathbb{E}[|X|^p])^{1/p}$.

According to [60, Proposition 5.10], we have the following concentration inequality for sub-Gaussian random variables.

Lemma 2. (Hoeffding-type inequality [60, Proposition 5.10]) Let $X_1, \ldots, X_N$ be independent zero-mean sub-Gaussian random variables, and let $K = \max_i \|X_i\|_{\psi_2}$. Then, for any $\alpha = [\alpha_1, \alpha_2, \ldots, \alpha_N]^T \in \mathbb{R}^N$ and any $t \geq 0$, it holds that

$$P\left( \left| \sum_{i=1}^N \alpha_i X_i \right| \geq t \right) \leq \exp\left( -\frac{ct^2}{K^2\|\alpha\|_2^2} \right), \quad (24)$$

where $c > 0$ is a constant.

Alongside the sub-Gaussian notion in Definition 1, we use the following definition of a sub-exponential random variable and sub-exponential norm.

Definition 2. A random variable $X$ is said to be sub-exponential if there exists a positive constant $C$ such that $(\mathbb{E}[|X|^p])^{1/p} \leq Cp$ for all $p \geq 1$. The sub-exponential norm of $X$ is defined as $\|X\|_{\psi_1} = \sup_{p \geq 1} p^{-1} (\mathbb{E}[|X|^p])^{1/p}$.

We have the following concentration inequality for sums of independent sub-exponential random variables.

Lemma 3. (Bernstein-type inequality [60, Proposition 5.16]) Let $X_1, \ldots, X_N$ be independent zero-mean sub-exponential random variables, and $K = \max_i \|X_i\|_{\psi_2}$. Then for every $\alpha = [\alpha_1, \ldots, \alpha_N]^T \in \mathbb{R}^N$ and $\epsilon \geq 0$, it holds that

$$P\left( \left| \sum_{i=1}^N \alpha_i X_i \right| \geq \epsilon \right) \leq 2 \exp\left( -c \cdot \min\left( \frac{\epsilon^2}{K^2\|\alpha\|_2^2}, \frac{\epsilon}{K\|\alpha\|_\infty} \right) \right), \quad (25)$$

where $c > 0$ is a constant.

B Proof of Theorem 1 (Recovery Guarantee for Amplitude-Based Loss Minimization)

Before proving the theorem, we present some additional auxiliary results.
B.1 Useful Lemmas

First, based on a basic two-sided concentration bound for standard Gaussian matrices (cf., Lemma 1 in Appendix A, and the well-established chaining arguments used in [5] [35], we have the following lemma, which essentially gives a two-sided Set-Restricted Eigenvalue Condition (S-REC) [5].

**Lemma 4.** [5 Lemma 4.1], [36 Lemma 2] For $\alpha < 1$ and $\delta > 0$, if $m = \Omega\left(\frac{k}{\alpha^2} \log \frac{Lr}{\delta}\right)$, then with probability $1 - e^{-\Omega(\alpha^2 m)}$, we have for all $x_1, x_2 \in G(B_2^k(r))$ that

\[
(1 - \alpha)\|x_1 - x_2\|_2 - \delta \leq \frac{1}{\sqrt{m}} \|A(x_1 - x_2)\|_2 \leq (1 + \alpha)\|x_1 - x_2\|_2 + \delta. \tag{26}
\]

In addition, we have the following lemma, which is similar to Lemma 1 in Appendix A.

**Lemma 5.** [61 Lemma 4.3] Suppose that $X_1, X_2, \ldots, X_N$ are i.i.d. standard normal random variables. Then, for $0 \leq \epsilon \leq \mu$, we have with probability $1 - e^{-\Omega(N\epsilon^2)}$ that

\[
\sqrt{N}(\mu - \epsilon) < \sqrt{\sum_{i=1}^{\lfloor N/2 \rfloor} |X_i|^2} < \sqrt{N}(\mu + \epsilon), \tag{27}
\]

where $\mu \geq \frac{1}{15} \sqrt{\frac{\pi}{2}}$ is a positive constant, and $|X|_1 \leq |X|_2 \leq \ldots \leq |X|_N$, i.e., $|X|_k$ is the $k$-th smallest entry in $\{X_1, X_2, \ldots, X_N\}$.

Based on Lemma 5 and using a chaining argument as in [5] [35], we arrive at the following lemma, which can be viewed as another (one-sided) variant of the S-REC in Lemma 4. The proof of Lemma 6 follows easily from [5] [35]: for completeness, we provide an outline in Appendix D.

**Lemma 6.** Let $\mu \geq \frac{1}{15} \sqrt{\frac{\pi}{2}}$ be the same positive constant as in Lemma 5. For $\alpha < \mu$ and $\delta > 0$, if $m = \Omega\left(\frac{k}{\alpha^2} \log \frac{Lr}{\delta}\right)$, then with probability $1 - e^{-\Omega(\alpha^2 m)}$, we have for all $x_1, x_2 \in G(B_2^k(r))$ that

\[
(\mu - \alpha)\|x_1 - x_2\|_2 - \delta \leq \min_{I \subseteq [m], |I| \geq \frac{1}{\pi}} \frac{1}{\sqrt{m}} \|A_I(x_1 - x_2)\|_2, \tag{28}
\]

where for any index set $I \subseteq [m]$, $A_I$ denotes the $|I| \times n$ sub-matrix of $A$ that only keeps the rows indexed by $I$.

B.2 Proof of Theorem 1

Since $p \in \text{Range}(G)$, we have

\[
\|y - |Ap|\|_2^2 + \tau \geq \min_{w \in \text{Range}(G)} \|y - |Aw|\|_2^2 + \tau \tag{29}
\]

\[
\geq \|y - |Aq|\|_2^2 \tag{30}
\]

\[
= \|y - |Ap| - (|Aq| - |Ap|)\|_2^2. \tag{31}
\]

where (30) follows from (5). Let $S, T$ be the following index sets:

\[
S := \{i \in [m] : \text{sign}(a_i^T p) = \text{sign}(a_i^T q)\}, \tag{32}
\]

\[
T := \{i \in [m] : \text{sign}(a_i^T p) \neq \text{sign}(a_i^T q)\}. \tag{33}
\]

Without loss of generality, we assume that $|S| \geq \frac{m}{2}$ (otherwise, we have $|T| \geq \frac{m}{2}$ and we can derive an upper bound for $\|q + x\|_2$ instead of for $\|q - x\|_2$). Expanding the squares in (31) and rearranging, we obtain

\[
\|Ap - |Aq|\|_2^2 \leq 2\|y - |Ap|, |Aq| - |Ap|\|_2^2 + \tau \tag{34}
\]

\[
\leq 2\|y - |Ap|\|_2^2\|Aq - |Ap|\|_2^2 + \tau \tag{35}
\]

\[
\leq 2\|y - |Ap|\|_2^2\|A(q - p)\|_2^2 + \tau \tag{36}
\]

\[
\leq 2\|\eta\|_2^2 + \|A(p - x)\|_2^2\|A(q - p)\|_2^2 + \tau, \tag{37}
\]

where (36) follows from the inequality $\|a - b\| \leq \min\{|a - b|, |a + b|\}$ for $a, b \in \mathbb{R}$, and (37) follows from (1) and the triangle inequality. Note that $p - x$ is a fixed vector. From a basic concentration
bound for standard Gaussian matrices (cf., Lemma 1 in Appendix A), we have with probability \(1 - e^{-\Omega(m)}\) that
\[
\|A(p - x)\|_2 \leq 2\sqrt{m}\|p - x\|_2.
\] (38)

In addition, setting \(\alpha = \frac{1}{2}\) in Lemma 4, we obtain that if \(m = \Omega\left(k \log \frac{1}{\delta}\right)\), then with probability \(1 - e^{-\Omega(m)}\), it holds that
\[
\|A(q - p)\|_2 \leq \sqrt{m} \left(\frac{3}{2}\|q - p\|_2 + \delta\right).
\] (39)

Moreover, setting \(\alpha = \frac{\mu}{2}\) in Lemma 6, we have that if \(m = \Omega\left(k \log \frac{1}{\delta}\right)\), with probability \(1 - e^{-\Omega(m)}\), it holds that
\[
\|Aq\|_2 - \|Ap\|_2 \leq \|A_S(p - q)\|_2 + \|A_T(p + q)\|_2 \geq \|A_S(p - q)\|_2 \geq m \left(\frac{\mu}{2}\|p - q\|_2 - \delta\right)^2 \geq m \left(\frac{\mu^2}{4}\|p - q\|_2^2 - \mu\delta\|p - q\|_2\right),
\] (40) (41) (42) (43)

where (40) follows from the definitions of \(S\) and \(T\), and (42) follows from Lemma 6 and our assumption \(|S| \geq \frac{m}{\eta}\). Combining (37), (38), (39), and (43), we obtain
\[
\frac{\mu^2}{4}\|p - q\|_2^2 - \mu\delta\|p - q\|_2 \leq 2 \left(\frac{\|\eta\|_2}{\sqrt{m}} + 2\|p - x\|_2\right) \left(\frac{3}{2}\|p - q\|_2 + \delta\right) + \tau \frac{m}{2},
\] (44)

which implies
\[
\|p - q\|_2 \leq O \left(\|p - x\|_2 + \frac{\|\eta\|_2}{\sqrt{m}} + \frac{\tau \frac{m}{2}}{\sqrt{m}}\right)\|p - x\|_2 + \left(\|p - x\|_2 + \frac{\|\eta\|_2}{\sqrt{m}}\right) \delta + O \left(\frac{\tau \frac{m}{2}}{\sqrt{m}}\right).
\] (45)

Simplifying terms in (45), we obtain
\[
\|p - q\|_2 \leq O \left(\|p - x\|_2 + \frac{\|\eta\|_2 + \sqrt{\tau}}{\sqrt{m}} + \delta\right).
\] (46)

From the triangle inequality \(\|q - x\|_2 \leq \|p - q\|_2 + \|p - x\|_2\), we obtain the desired result.

C Proof of Theorem 2 (Spectral Initializations with Generative Priors)

Before proving Theorem 2, we provide a simplified outline and some useful auxiliary results. In this appendix, it is particularly important to remember that \(c\) and \(c'\) represent small constants whose values may differ from line to line.

C.1 Simplified Outline of the Proof

Since the full proof of Theorem 2 is rather lengthy, we first provide an outline with certain simplifying assumptions (which are non-rigorous but will be dropped subsequently). Specifically, we consider an idealized scenario in which \(\eta = 0\) (i.e., the noiseless setting), and \(\lambda\) exactly equals its expectation \(\|x\|_2\). Then, we have \(y_i = |\langle a_i, x \rangle|\), and \(V\) simplifies to
\[
V_0 = \frac{1}{m} \sum_{i=1}^m y_i a_i a_i^T 1_{\{\langle a_i^T \hat{x} \rangle < \mu\}}.
\] (47)

In this simplified scenario, the mean is given by (cf. (13))
\[
J = \mathbb{E}[V_0] = \|x\|_2(\gamma_0 I_n + \beta_0 \hat{x} \hat{x}^T),
\] (48)

and \(\hat{x} \in S^{n-1}\) is the vector that satisfies
\[
\hat{x} = \arg\max_{w \in G(D)} w^T V_0 w.
\] (49)
On the one hand, since $\bar{x} \in \hat{G}(\mathcal{D})$, we have
\[
\hat{x}^T V_0 \hat{x} \geq \bar{x}^T V_0 \bar{x}
\]
where (52) uses (48). With $y_i = |a_i^T x|$, we have that $y_i (a_i^T \hat{x})^2 1_{\{t < |a_i^T \bar{x}| < u \}}$ is upper bounded by $\|x\|_2 u^3$. From (52) we can use the concentration inequality for the sum of sub-Gaussian random variables (cf. Lemma 2) to derive a lower bound of $\hat{x}^T V_0 \hat{x}$.

On the other hand, we have
\[
\hat{x}^T V_0 \hat{x} = \hat{x}^T (V_0 - J) \hat{x} + \hat{x}^T J \hat{x}
\]
where we again used (48). Note that $\hat{x}$ is dependent on $V_0$. To upper bound (54), we construct a $\delta$-net, and write $\hat{x}$ as
\[
\hat{x} = (\hat{x} - \bar{x}) + \bar{x},
\]
where $\bar{x}$ is in the $\delta$-net satisfying $\|\hat{x} - \bar{x}\|_2 < \delta$. For any $\bar{s} \in \mathbb{R}^n$, we have that $y_i (a_i^T \bar{s})^2 1_{\{t < |a_i^T \bar{x}| < u \}}$ is sub-exponential, with the sub-exponential norm being upper bounded by $\|x\|_2 u C$, where $C$ is an absolute constant. Using a concentration inequality for the sum of sub-exponential random variables (cf. Lemma 3) and taking a union bound over the $\delta$-net, we obtain an upper bound for $\hat{x}^T (V_0 - J) \hat{x}$. Using a high-probability upper bound on $\|a_i\|$ and the fact that $\|\hat{x} - \bar{x}\|_2 < \delta$, we can control the terms $\hat{x}^T (V_0 - J) \hat{x}$ and $\hat{x}^T (V_0 - J) (\hat{x} - \bar{x})$. Then, from (54), we obtain an upper bound on $\hat{x}^T V_0 \hat{x}$. Combining the upper and lower bounds on $\hat{x}^T V_0 \hat{x}$, we can derive the desired result.

In the full analysis in Section C.3 we additionally carefully deal with the noise terms and the case that $\lambda$ approximately equals $\|x\|_2$.

### C.2 Additional Lemmas

Based on Lemma 2 we present the following simple lemma showing that $\lambda$ approximates $\|x\|_2$.

**Lemma 7.** For any fixed $c \in (0, 1)$, with probability $1 - e^{-\Omega(m)}$, we have that $\lambda$ defined in (11) satisfies
\[
1 - c < \frac{\lambda}{\|x\|_2} < 1 + c.
\]

**Proof.** Since $|\langle a_i, x \rangle|$ is sub-Gaussian with the sub-Gaussian norm upper bounded by $C \|x\|_2$ and $\mathbb{E}[|\langle a_i, x \rangle|] = \sqrt{\frac{2}{\pi}} \|x\|_2$, from Lemma 2 we have with probability $1 - e^{-\Omega(m)}$ that
\[
\left| \frac{1}{m} \sum_{i=1}^{m} |\langle a_i, x \rangle| - \sqrt{\frac{2}{\pi}} \|x\|_2 \right| \leq c \|x\|_2.
\]

In addition, by the Cauchy–Schwarz inequality and the upper bound for the noise term assumed in (8), we have
\[
\left| \frac{1}{m} \sum_{i=1}^{m} \eta_i \right| \leq \frac{\|\eta\|_2}{\sqrt{m}} \leq c_0 \|x\|_2.
\]

Then, using (57), (58), and the triangle inequality, we obtain
\[
\left( \sqrt{\frac{2}{\pi}} - c' \right) \|x\|_2 \leq \frac{1}{m} \sum_{i=1}^{m} y_i = \frac{1}{m} \sum_{i=1}^{m} (|\langle a_i, x \rangle| + \eta_i) \leq \left( \sqrt{\frac{2}{\pi}} + c' \right) \|x\|_2.
\]

Therefore, we obtain
\[
(1 - c) \|x\|_2 < \lambda = \sqrt{\frac{2}{\pi}} \left( \frac{1}{m} \sum_{i=1}^{m} y_i \right) < (1 + c) \|x\|_2
\]
for a choice of $c$ possibly different from that above, but still arbitrarily small. 

\[\square\]
Next, we present the following lemma concerning the expectation of the truncated empirical matrix.

**Lemma 8.** Suppose that $a \sim \mathcal{N}(0, I_n)$. For any $s \in \mathbb{R}^n$ and $\eta \in \mathbb{R}$, choosing $\alpha_1, \alpha_2 \in \mathbb{R}$ such that $1 + \frac{\eta}{\|s\|^2} < \alpha_1 < \alpha_2$ and setting $y = |\langle a, s \rangle| + \eta$, we have

$$
E \left[ y a a^T 1_{\{\alpha_1 \|s\|^2 < y < \alpha_2 \|s\|^2 \}} \right] = \|s\|^2 (\gamma_1 I_n + \beta_1 \bar{s}s^T) + \eta (\gamma_1 I_n + \bar{\beta}_1 \bar{s}s^T),
$$

(61)

where $\bar{s} := \frac{s}{\|s\|^2}$, and for $g \sim \mathcal{N}(0, 1)$ and $\phi_1(x) = 1_{\{a_1 - \frac{\eta}{\|a\|^2} < |x| < a_2 - \frac{\eta}{\|a\|^2}\}}$, we define $\gamma_1 = E[|g|\phi_1(g)]$, $\beta_1 = E[|g|^3\phi_1(g)] - \gamma_1$, $\bar{\gamma}_1 = E[\phi_1(g)]$ and $\bar{\beta}_1 = E[|g|^2\phi_1(g)] - \bar{\gamma}_1$. (Note that from the assumption that $1 + \frac{\eta}{\|s\|^2} < \alpha_1 < \alpha_2$, we have $\beta_1 > 0$ and $\bar{\beta}_1 > 0$.)

**Proof.** Let $g = \langle a, \bar{s} \rangle \sim \mathcal{N}(0, 1)$. We have

$$
y a a^T 1_{\{\alpha_1 \|s\|^2 < y < \alpha_2 \|s\|^2 \}} = (\|s\|^2 |a^T \bar{s}| + \eta) a a^T 1_{\{\alpha_1 - \frac{\eta}{\|a\|^2} < |\langle a, s \rangle| < \alpha_2 - \frac{\eta}{\|a\|^2}\}}
$$

(62)

$$
= (\|s\|^2 |g| + \eta) \phi_1(g) a a^T.
$$

(63)

For any $i \in [n]$, we have $\text{Cov}[a_i, g] = \bar{s}_i$, and thus $a_i$ can be written as

$$
a_i = \bar{s}_i g + t_i,
$$

(64)

where $t_i \sim \mathcal{N}(0, 1 - \bar{s}_i^2)$ is independent of $g$. Then, we have

$$
E \left[ |g|\phi_1(g)a_i^2 \right] = E \left[ |g|\phi_1(g)(\bar{s}_i g + t_i)^2 \right] = \bar{s}_i^2 E \left[ |g|^2\phi_1(g) \right] + (1 - \bar{s}_i^2) E[|g|\phi_1(g)]
$$

(65)

$$
= \gamma_1 + \beta_1 \bar{s}_i^2,
$$

(66)

and similarly,

$$
E \left[ \phi_1(g)a_i^2 \right] = E \left[ \phi_1(g)(\bar{s}_i g + t_i)^2 \right] = \bar{s}_i^2 E \left[ g^2\phi_1(g) \right] + (1 - \bar{s}_i^2) E[\phi_1(g)]
$$

(67)

$$
= \bar{\gamma}_1 + \bar{\beta}_1 \bar{s}_i^2.
$$

(68)

Moreover, for $1 \leq i \neq j \leq n$, we have $0 = E[a_ia_j] = E[(\bar{s}_i g + t_i)(\bar{s}_j g + t_j)] = \bar{s}_i \bar{s}_j + E[t_it_j]$, which gives $E[t_it_j] = -\bar{s}_i \bar{s}_j$. Then, similarly to (65) and (68), we have

$$
E \left[ |g|\phi_1(g)a_ia_j \right] = E \left[ |g|\phi_1(g)(\bar{s}_i g + t_i)(\bar{s}_j g + t_j) \right] = \beta_1 \bar{s}_i \bar{s}_j,
$$

(69)

and

$$
E \left[ \phi_1(g)a_ia_j \right] = E \left[ \phi_1(g)(\bar{s}_i g + t_i)(\bar{s}_j g + t_j) \right] = \bar{\beta}_1 \bar{s}_i \bar{s}_j.
$$

(70)

Combining (63), (67), (70), (69), (71), we obtain the desired result. □

### C.3 Proof of Theorem 2

Let $\mathcal{E}$ be the event that

$$1 - c \leq \frac{\lambda}{\|x\|^2} \leq 1 + c.
$$

(73)

From Lemma 7, we know that $\mathcal{E}$ occurs with probability $1 - e^{-O(n)}$. Throughout the following, we assume that $\mathcal{E}$ holds, and the relevant probabilities are conditioned accordingly.
Lower bounding $\hat{x}^T \hat{V} \hat{x}$: Since $\hat{x} \in \hat{G}(D)$, and $\hat{x}$ is a solution of (14), we have

$$\hat{x}^T \hat{V} \hat{x} \geq \hat{x}^T \hat{V} \hat{x}$$

$$= \frac{1}{m} \sum_{i=1}^{m} y_i \left( a_i^T \hat{x} \right)^2 1_{\{\lambda < y_i < \mu \lambda \}}$$

$$= \frac{1}{m} \sum_{i=1}^{m} |a_i^T \hat{x}|^3 1_{\{\lambda < y_i < \mu \lambda \}} + \frac{1}{m} \sum_{i=1}^{m} \eta_i \left( a_i^T \hat{x} \right)^2 1_{\{\lambda < y_i < \mu \lambda \}}$$

$$\geq \frac{1}{m} \sum_{i=1}^{m} |a_i^T \hat{x}|^3 1_{\{(1+\epsilon')\|\hat{x}\| < \mu(1+\epsilon')\|\hat{x}\| \}} - \frac{1}{m} \sum_{i=1}^{m} \eta_i \left( a_i^T \hat{x} \right)^2 1_{\{(1-\epsilon')\|\hat{x}\| < \lambda(1-\epsilon')\|\hat{x}\| \}}$$

$$= \frac{1}{m} \sum_{i=1}^{m} |a_i^T \hat{x}|^3 1_{\{(1+\epsilon')|\|a_i^T \hat{x}\| < \mu(1+\epsilon')\|a_i^T \hat{x}\| \}} - \frac{1}{m} \sum_{i=1}^{m} \eta_i \left( a_i^T \hat{x} \right)^2 1_{\{(1-\epsilon')|\|a_i^T \hat{x}\| < \lambda(1-\epsilon')\|a_i^T \hat{x}\| \}}$$

$$\geq \frac{1}{m} \sum_{i=1}^{m} |a_i^T \hat{x}|^3 1_{\{(1+\epsilon')|\|a_i^T \hat{x}\| < \mu(1+\epsilon')\|a_i^T \hat{x}\| \}} - \frac{1}{m} \sum_{i=1}^{m} \eta_i \left( a_i^T \hat{x} \right)^2 1_{\{(1-\epsilon')|\|a_i^T \hat{x}\| < \lambda(1-\epsilon')\|a_i^T \hat{x}\| \}}$$

where we use (73) in (77), we use $y_i = |a_i^T \hat{x}| + \eta_i$ in (76) and (78), and we use (9) as well as $u > l > 1 + c_1$ to derive (79). We aim to bound the two terms in (79). Let

$$U = \frac{1}{m} \sum_{i=1}^{m} |a_i^T \hat{x}| a_i a_i^T 1_{\{(1+\epsilon')|\|a_i^T \hat{x}\| < \mu(1+\epsilon')\|a_i^T \hat{x}\| \}}$$

$$W = \frac{1}{m} \sum_{i=1}^{m} |a_i^T \hat{x}| a_i a_i^T 1_{\{(1-\epsilon')|\|a_i^T \hat{x}\| < \lambda(1-\epsilon')\|a_i^T \hat{x}\| \}}$$

$$\hat{U} = \frac{1}{m} \sum_{i=1}^{m} \eta_i |a_i| a_i a_i^T 1_{\{(1+\epsilon')|\|a_i^T \hat{x}\| < \mu(1+\epsilon')\|a_i^T \hat{x}\| \}}$$

$$\hat{W} = \frac{1}{m} \sum_{i=1}^{m} |a_i| a_i a_i^T 1_{\{(1-\epsilon')|\|a_i^T \hat{x}\| < \lambda(1-\epsilon')\|a_i^T \hat{x}\| \}}$$

According to the proof of Lemma 8 we have

$$E[U] = \|x\|_2 \left( \gamma I_n + \beta x^T x \right)$$

$$E[W] = \|x\|_2 \left( \gamma' I_n + \beta' x^T x \right)$$

where for $g \sim \mathcal{N}(0, 1)$ and

$$\phi(x) := 1_{\{(1+\epsilon')|\|x\| < \mu(1+\epsilon')\|x\| \}}$$

$$\psi(x) := 1_{\{(1-\epsilon')|\|x\| < \lambda(1-\epsilon')\|x\| \}}$$

we set $\gamma = E[|g|\phi(g)]$, $\beta = E[|g|^3\phi(g)] - \gamma$, $\gamma' = E[|g|\psi(g)]$, $\beta' = E[|g|^3\psi(g)] - \gamma'$, $\bar{\eta} = \frac{1}{m} \sum_{i=1}^{m} \eta_i$, $\bar{\gamma} = E[\phi(g)]$, $\bar{\beta} = E[g^2\phi(g)] - \bar{\gamma}$, $\bar{\gamma}' = E[\psi(g)]$, and $\bar{\beta}' = E[g^2\psi(g)] - \bar{\gamma}'$. Then, in (79), we have

$$\|x\|_2 \sum_{i=1}^{m} |a_i^T \hat{x}|^3 1_{\{(1+\epsilon')|\|a_i^T \hat{x}\| < \mu(1+\epsilon')\|a_i^T \hat{x}\| \}}$$

$$= \hat{x}^T U \hat{x}$$
\begin{align*}
= \hat{x}^T(U - E[U])\hat{x} + \hat{x}^TE[U]\hat{x} & \quad (91) \\
= \frac{1}{m} \sum_{i=1}^m (a_i^T \hat{x})(a_i^T \hat{x})^2 \phi(a_i^T \hat{x}) - E \left[ (a_i^T \hat{x})(a_i^T \hat{x})^2 \phi(a_i^T \hat{x}) \right] + \|x\|_2(\gamma + \beta), & \quad (92)
\end{align*}

where we use (84), (88), and \(\|x\|_2 = 1\) to obtain (92). Note that \(|a_i^T \hat{x}|(a_i^T \hat{x})^2 \phi(a_i^T \hat{x}) = |a_i^T \hat{x}|(a_i^T \hat{x})^2 \phi(a_i^T \hat{x}) \leq u(1 - c)^3 \|x\|_2 = O(\|x\|_2)\). Then, from Lemma 2, we have with probability \(1 - e^{-\Omega(m)}\) that

\begin{equation}
\left| \frac{1}{m} \sum_{i=1}^m (a_i^T \hat{x})(a_i^T \hat{x})^2 \phi(a_i^T \hat{x}) - E \left[ (a_i^T \hat{x})(a_i^T \hat{x})^2 \phi(a_i^T \hat{x}) \right] \right| \leq c\|x\|_2. \quad (93)
\end{equation}

In addition, we have

\begin{equation}
\frac{1}{m} \sum_{i=1}^m |\eta_i| (a_i^T \hat{x})^2 1_{\{(1-c)\|x\|_2 < \|x\|_2(1+c)\}} \leq \hat{x}^T \hat{\mathbf{W}} \hat{x} \quad (94)
\end{equation}

\begin{equation}
\hat{x}^T(W - E[W])\hat{x} + \hat{x}^TE[W]\hat{x} \leq \frac{1}{m} \sum_{i=1}^m |\eta_i| \left( (a_i^T \hat{x})^2 \psi(a_i^T \hat{x}) - E \left[ (a_i^T \hat{x})^2 \psi(a_i^T \hat{x}) \right] \right) + \eta(\gamma' + \beta') \quad (95)
\end{equation}

\begin{equation}
\leq \frac{1}{m} \sum_{i=1}^m |\eta_i| \left( (a_i^T \hat{x})^2 \psi(a_i^T \hat{x}) - E \left[ (a_i^T \hat{x})^2 \psi(a_i^T \hat{x}) \right] \right) + c_0\|x\|_2(\gamma' + \beta'), \quad (96)
\end{equation}

where we use \(\hat{\eta} = \frac{1}{m} \sum_{i=1}^m |\eta_i| \leq \frac{\|\eta\|_2}{\sqrt{m}} \leq c_0\|x\|_2\) (see (8)) in (97). From the definition of \(\psi\) in (84), we have \((a_i^T \hat{x})^2 \psi(a_i^T \hat{x}) \leq u^2(1+c)^2 = O(1)\). Setting \(\alpha_i = \frac{|m|}{m}\) and \(X_i = (a_i^T \hat{x})^2 \psi(a_i^T \hat{x}) - E[(a_i^T \hat{x})^2 \psi(a_i^T \hat{x})]\) in Lemma 2, we obtain with probability \(1 - e^{-\Omega(m)}\) that

\begin{equation}
\left| \frac{1}{m} \sum_{i=1}^m |\eta_i| \left( (a_i^T \hat{x})^2 \psi(a_i^T \hat{x}) - E \left[ (a_i^T \hat{x})^2 \psi(a_i^T \hat{x}) \right] \right) \right| \leq \frac{\|\eta\|_2}{\sqrt{m}} \leq c_0\|x\|_2. \quad (98)
\end{equation}

Combining (97), (92), (93), (96) and (98), we have with probability \(1 - e^{-\Omega(m)}\) that

\begin{equation}
\hat{x}^T V \hat{x} \geq (\gamma + \beta - c)\|x\|_2. \quad (99)
\end{equation}

**Upper bounding \(\hat{x}^T V \hat{x}\):** We have

\begin{align*}
\hat{x}^T V \hat{x} &= \frac{1}{m} \sum_{i=1}^m |a_i^T \hat{x}|(a_i^T \hat{x})^2 1_{\{(1-c)<|a_i^T \hat{x}|<u(1+c)\}} + \frac{1}{m} \sum_{i=1}^m |\eta_i|(a_i^T \hat{x})^2 1_{\{(1-c)<|a_i^T \hat{x}|<u(1+c)\}} \quad (100) \\
&\leq \frac{1}{m} \sum_{i=1}^m |a_i^T \hat{x}|(a_i^T \hat{x})^2 1_{\{(1-c)<|a_i^T \hat{x}|<u(1+c)\}} + \frac{1}{m} \sum_{i=1}^m |\eta_i|(a_i^T \hat{x})^2 1_{\{(1-c)<|a_i^T \hat{x}|<u(1+c)\}} \quad (101)
\end{align*}

\begin{equation}
= \hat{x}^T \hat{\mathbf{W}} \hat{x} + \hat{x}^T \hat{\mathbf{W}} \hat{x} \quad (102)
\end{equation}

\begin{equation}
= \hat{x}^T(W - E[W])\hat{x} + \hat{x}^T(W - E[W])\hat{x} + \|x\|_2(\gamma' + \beta' (\hat{x}^T \hat{x})) + \eta(\gamma' + \beta' (\hat{x}^T \hat{x})) \quad (103)
\end{equation}

\begin{equation}
\leq \hat{x}^T(W - E[W])\hat{x} + \hat{x}^T(W - E[W])\hat{x} + \|x\|_2(\gamma' + \beta' (\hat{x}^T \hat{x})) + c\|x\|_2, \quad (104)
\end{equation}

where (101) is derived similarly to (79), (102) comes from the definitions of \(\mathbf{W}\) and \(\hat{\mathbf{W}}\) in (81) and (83). (103) follows from (85) and (87), and (104) is deduced from \(\hat{\eta} \leq c_0\|x\|_2\).

We aim to provide an upper bound for (104). For \(\delta > 0\), let \(M\) be a \((\delta/L)-\)net of \(D \subseteq B_2^L(r)\). There exists such a net with [60, Lemma 5.2]

\begin{equation}
\log |M| \leq k \log \frac{4Lr}{\delta}. \quad (105)
\end{equation}
By the $\tilde{L}$-Lipschitz continuity of $\tilde{G}$, we have that $\tilde{G}(M)$ is a $\delta$-net of $\tilde{G}(D)$. We write

$$\tilde{x} = (\hat{x} - \bar{x}) + \hat{x},$$  \hspace{1cm} (106)

where $\tilde{x} \in \tilde{G}(M)$ with $\|\tilde{x} - \bar{x}\|_2 \leq \delta$. Then, we have

$$\tilde{x}^T(W - \mathbb{E}[W])\tilde{x} = \hat{x}^T(W - \mathbb{E}[W])\hat{x} + 2(\hat{x} - \tilde{x})^T(W - \mathbb{E}[W])\tilde{x} + (\hat{x} - \tilde{x})^T(W - \mathbb{E}[W])(\hat{x} - \tilde{x}).$$  \hspace{1cm} (107)

We bound the three terms in (107) separately.

1. **The first term:** For any $s \in \tilde{G}(M)$, we have from (81) that

$$s^T(W - \mathbb{E}[W])s = \frac{1}{m} \sum_{i=1}^{m} (|a_i^Tx||a_i|)^2 1_{\{|l(1-c)|<|a_i^T\bar{x}|<u(1+c)|\}} - s^T\mathbb{E}[W]s.$$  \hspace{1cm} (108)

Note that for $s \in \tilde{G}(M)$, $\|s\|_2 = 1$. Then, $|a_i^Tx||a_i|s^2 1_{\{|l(1-c)|<|a_i^T\bar{x}|<u(1+c)|\}}$ is sub-exponential with the sub-exponential norm upper bounded by $Cu(1+c)\|x\|_2$. From Lemma 3 we have with probability at least $1 - e^{-\Omega(m^2)}$

$$|s^T(W - \mathbb{E}[W])s| \leq \epsilon \|x\|_2.$$  \hspace{1cm} (109)

Taking a union bound over all $\tilde{G}(M)$, we obtain that when $m = \Omega(k \log \frac{Lr}{\epsilon})$, with probability $1 - e^{-\Omega(m^2)}$, for all $s \in \tilde{G}(M)$,

$$|s^T(W - \mathbb{E}[W])s| \leq \epsilon \|x\|_2.$$  \hspace{1cm} (110)

Hence, since $\bar{x} \in \tilde{G}(M)$, setting $\epsilon$ to be a sufficiently small absolute constant, we obtain that when $m = \Omega(k \log \frac{Lr}{\epsilon})$, with probability $1 - e^{-\Omega(m)}$,

$$|\tilde{x}^T(W - \mathbb{E}[W])\tilde{x}| \leq \epsilon \|x\|_2.$$  \hspace{1cm} (111)

2. **The second term:** From Lemma 2 and a union bound over $[m]$, we have with probability $1 - me^{-\Omega(n)}$ that

$$\max_{i \in [m]} \|a_i\|_2 \leq \sqrt{2n}.$$  \hspace{1cm} (112)

Conditioned on the event in (112), we have

$$(\hat{x} - \bar{x})^T(W - \mathbb{E}[W])\hat{x}$$

$$= \frac{1}{m} \sum_{i=1}^{m} (|a_i^T\tilde{x}|(|a_i^T\tilde{x}|)(|a_i^T\hat{x}|)(|a_i^T\hat{x}|)1_{\{|l(1-c)|<|a_i^T\tilde{x}|<u(1+c)|\}} - (\hat{x} - \tilde{x})^T\mathbb{E}[W]\hat{x})$$

$$= \frac{1}{m} \sum_{i=1}^{m} |a_i^T\tilde{x}|(|a_i^T\tilde{x}|)(|a_i^T\hat{x}|)(|a_i^T\hat{x}|)1_{\{|l(1-c)|<|a_i^T\tilde{x}|<u(1+c)|\}}$$

$$- \|x\|_2 \left(\gamma'(\hat{x} - \tilde{x})^T\hat{x} + \beta'(\hat{x} - \tilde{x})^T\tilde{x}\right)$$

$$\leq (2n(1+c)u + \gamma' + \beta')\|x\|_2\delta,$$  \hspace{1cm} (114)

where we use (112), $\|\hat{x} - \bar{x}\|_2 \leq \delta$, $\|x\|_2 = 1$, and the standard inequality $|a^Tb| \leq \|a\|_2 \|b\|_2$ (for any $a$ and $b$) in (115). Therefore, we obtain with probability $1 - e^{-\Omega(n)}$ that

$$|2(\hat{x} - \bar{x})^T(W - \mathbb{E}[W])\tilde{x}| \leq 2(2n(1+c)u + \gamma' + \beta')\|x\|_2\delta.$$  \hspace{1cm} (116)

3. **The third term:** By an argument similar to (116), we have with probability $1 - e^{-\Omega(n)}$ that

$$|\hat{x}^T(W - \mathbb{E}[W])\tilde{x}| \leq (2n(1+c)u + \gamma' + \beta')\|x\|_2\delta.$$  \hspace{1cm} (117)

Combining (107), (111), (116), and (117), and setting $\delta = \frac{\epsilon}{n}$, we obtain that when $m = \Omega(k \log(Lnr))$, with probability $1 - e^{-\Omega(m)}$, it holds that

$$|\tilde{x}^T(W - \mathbb{E}[W])\tilde{x}| \leq \epsilon \|x\|_2.$$  \hspace{1cm} (118)
Next, we provide an upper bound for $\hat{x}^T(\tilde{W} - \mathbb{E}[W])\hat{x}$. The strategy is mostly similar to that for upper bounding $\hat{x}^T(\tilde{W} - \mathbb{E}[W])\hat{x}$ in (104), but we provide the details below for completeness. Similarly to (107), we have

$$\hat{x}^T(\tilde{W} - \mathbb{E}[W])\hat{x} = \hat{x}^T(\tilde{W} - \mathbb{E}[W])\hat{x} + 2(\hat{x} - \tilde{x})^T(\tilde{W} - \mathbb{E}[W])(\hat{x} - \tilde{x}).$$

(119)

We bound the three terms in (119) separately.

- **The first term:** From the definition of $W$ in (83), we have for any $s \in \hat{G}(M)$ that

$$s^T(\tilde{W} - \mathbb{E}[W])s = \frac{1}{m} \sum_{i=1}^{m} |\eta_i| (a_i^T s)^2 1_{|\langle a_i^T \hat{x} \rangle| < u(1 + \epsilon)} - s^T \mathbb{E}[W]s. \quad (120)$$

Note that $\|s\|_2 = 1$. For any $i \in [m]$, $(a_i^T s)^2 1_{|\langle a_i^T \hat{x} \rangle| < u(1 + \epsilon)}$ is sub-exponential with the sub-exponential norm upper bounded by an absolute constant $C$. Then, from Lemma 3 (see also [21, Theorem 3.1]), we obtain that for $t > 2$, with probability $1 - e^{-t}$, it holds that

$$|s^T(\tilde{W} - \mathbb{E}[W])s| \leq C \left( \frac{\|\eta\|_2}{m} \sqrt{t} + \frac{\|\eta\|_{\infty}}{m} t \right). \quad (121)$$

Setting $t = m$ in (121), and from $\|\eta\|_2 \leq c_0 \|x\|_2$ (cf. (8)) and $\|\eta\|_{\infty} \leq c_1 \|x\|_2$ (cf. (9)), we obtain with probability $1 - e^{-m}$ that

$$|s^T(\tilde{W} - \mathbb{E}[W])s| \leq c \|x\|_2. \quad (122)$$

Taking a union bound over all $\hat{G}(M)$, we obtain that when $m = \Omega(k \log \frac{Lr}{\delta})$, with probability $1 - e^{-\Omega(m^2)}$, for all $s \in \hat{G}(M)$,

$$|s^T(\tilde{W} - \mathbb{E}[W])s| \leq c \|x\|_2. \quad (123)$$

Hence, since $\hat{x} \in \hat{G}(M)$, we have

$$|\hat{x}^T(\tilde{W} - \mathbb{E}[W])\hat{x}| \leq c \|x\|_2. \quad (124)$$

- **The second term:** Conditioned on the event in (112), we have

$$\begin{align*}
(\hat{x} - \tilde{x})^T(\tilde{W} - \mathbb{E}[W])\tilde{x} &= \frac{1}{m} \sum_{i=1}^{m} |\eta_i| (a_i^T (\hat{x} - \tilde{x})) (a_i^T \tilde{x}) 1_{|\langle a_i^T \hat{x} \rangle| < u(1 + \epsilon)} \\
&\quad - \bar{\eta} (\hat{\gamma}'(\hat{x} - \tilde{x})^T \hat{x} + \hat{\beta}'((\hat{x} - \tilde{x})^T \tilde{x})) \\
&\leq \bar{\eta} (2n + \hat{\gamma}' + \hat{\beta}'). \quad (125)
\end{align*}$$

Therefore, we obtain with probability $1 - e^{-\Omega(n)}$ that

$$|2(\hat{x} - \tilde{x})^T(\tilde{W} - \mathbb{E}[W])\tilde{x}| \leq 2\bar{\eta} (2n + \hat{\gamma}' + \hat{\beta}'). \quad (127)$$

- **The third term:** By an argument similar to (127), we have with probability $1 - e^{-\Omega(n)}$ that

$$\begin{align*}
\left| (\hat{x} - \tilde{x})^T(\tilde{W} - \mathbb{E}[W])((\hat{x} - \tilde{x}) \tilde{x}) \right| &\leq (2n(1 + c)u + \gamma' + \beta')\bar{\eta}\delta^2. \quad (128)
\end{align*}$$

Note that $\bar{\eta} \leq \frac{\|\eta\|_2}{\sqrt{m}} \leq c_0 \|x\|_2$ (cf. (8)). Combining (119), (124), (127), and (128), and setting $\delta = \frac{\epsilon}{n}$, we obtain that when $m = \Omega(k \log(\hat{Lnr}))$, with probability $1 - e^{-\Omega(m)}$, it holds that

$$|\hat{x}^T(\tilde{W} - \mathbb{E}[W])\tilde{x}| \leq c \|x\|_2. \quad (129)$$

Combining (104), (118), and (129), we obtain that when $m = \Omega(k \log(\hat{Lnr}))$, with probability $1 - e^{-\Omega(m)}$, it holds that

$$\hat{x}^TV\hat{x} \leq (c + \gamma' + \beta'(\hat{x}^T \hat{x})^2) \|x\|_2. \quad (130)$$

22
Combining and simplifying: Finally, combining (99) and (130), we obtain that when $m = \Omega(k \log(Lnr))$, with probability $1 - e^{-\Omega(m)}$,

$$
(\gamma + \beta - c)\|x\|_2 \leq (c + \gamma' + \beta'(\hat{x}^T \hat{x})^2) \|x\|_2.
$$

Simplifying (131), we obtain

$$
\beta'(\hat{x}^T \hat{x})^2 \geq \beta + (\gamma - \gamma') - 2c.
$$

Defining $\Phi(x) = \int_{-\infty}^{x} \frac{1}{\sqrt{2\pi}} e^{-\frac{t^2}{2}} dt$ to be the standard normal distribution function and setting $C_0 = \max_{x \in \mathbb{R}} \Phi'(x) = \frac{1}{\sqrt{2\pi}}$, we have for $g \sim \mathcal{N}(0, 1)$ that

$$
\gamma' - \gamma = E[|g|1_{\{t(1-c) \leq |g| \leq t(1+c)\}}] + E[|g|1_{\{u(1-c) \leq |g| \leq u(1+c)\}}]
$$

$$
\leq 2E[|g|1_{\{u(1-c) \leq |g| \leq u(1+c)\}}]
$$

$$
\leq 4u(1+c)P(u(1-c) < g < u(1+c))
$$

$$
\leq 5u(\Phi(u(1+c)) - \Phi(u(1-c)))
$$

$$
\leq 10C_0 u^2 c = O(c),
$$

where (134) uses the definitions of $\gamma$ and $\gamma'$ following (89). (135) multiplies by two due to the replacement of $|g|$ by $g$ in the probability, and (136) holds for small enough $c$. Similarly, we have $\beta \geq \beta' - O(c)$. Recall from (73) that the event $\mathcal{E}$ occurs with probability $1 - e^{-\Omega(m)}$. Therefore, from (132), we obtain that when $m = \Omega(k \log(Lnr))$, with probability $1 - e^{-\Omega(m)}$, it holds that

$$
(\hat{x}^T \hat{x})^2 \geq 1 - c'.
$$

Without loss of generality, we assume that $\hat{x}^T \hat{x} > 0$ (otherwise, we can similarly derive an upper bound for $\|\hat{x} + \hat{x}\|_2$ instead of $\|\hat{x} - \hat{x}\|_2$). Then, we have

$$
\|\hat{x} - \hat{x}\|_2^2 = 2(1 - \hat{x}^T \hat{x}) \leq 2(1 - (\hat{x}^T \hat{x})^2) \leq 2c'.
$$

By suitably renaming the constant, we obtain (16) as desired. Using Lemma 7, we also obtain (17).

D Proof of Lemma 6 (Variant of the S-REC)

For fixed $\delta > 0$ and a positive integer $\ell$, let $M = M_0 \subseteq M_1 \subseteq \ldots \subseteq M_\ell$ be a chain of nets of $B^2_2(r)$ such that $M_i$ is a $\delta_i$-net with $\delta_i = \frac{\delta}{2^i}$. There exists such a chain of nets with [60, Lemma 5.2]

$$
\log |M_i| \leq k \log \frac{4Lr}{\delta_i}.
$$

By the $L$-Lipschitz assumption on $G$, we have for any $i \in [\ell]$ that $G(M_i)$ is a $\delta_i$-net of $G(B^2_2(r))$.

We write

$$
x_1 = (x_1 - s_\ell) + \sum_{i=1}^{\ell} (s_i - s_{i-1}) + s_0,
$$

$$
x_2 = (x_2 - t_\ell) + \sum_{i=1}^{\ell} (t_i - t_{i-1}) + t_0,
$$

where $s_i, t_i \in G(M_i)$ for all $i \in [\ell]$, and $\|x_1 - s_\ell\| \leq \frac{\delta}{2^\ell}, \|x_2 - t_\ell\| \leq \frac{\delta}{2^\ell}, \|s_i - s_{i-1}\| \leq \frac{\delta}{2^\ell}, \|t_i - t_{i-1}\| \leq \frac{\delta}{2^\ell}$, and $\|s_i - s_{i-1}\| \leq \frac{\delta}{2^\ell}$ for all $i \in [\ell]$. Therefore, the triangle inequality gives

$$
\|x_1 - s_0\| \leq 2\delta, \quad \|x_2 - t_0\| \leq 2\delta.
$$

For any index $I \subseteq [m]$ with $|I| \geq \frac{m}{2}$, from the triangle inequality, we have

$$
\|A_I(x_1 - x_2)\|_2 \geq \|A_I(s_0 - t_0)\|_2 - \sum_{i=1}^{\ell} (\|A_I(s_i - s_{i-1})\|_2 + \|A_I(t_i - t_{i-1})\|_2)
$$

$$
- \|A_I(s_\ell - x_1)\|_2 - \|A_I(t_\ell - x_2)\|_2
$$

$$
\geq \|A_I(s_0 - t_0)\|_2 - \sum_{i=1}^{\ell} (\|A(s_i - s_{i-1})\|_2 + \|A(t_i - t_{i-1})\|_2)
$$

$$
- \|A(s_\ell - x_1)\|_2 - \|A(t_\ell - x_2)\|_2.
$$
For any \( r_1, r_2 \in G(M) \) and \( 0 < \alpha < \mu \), from Lemma 5 we have with probability \( 1 - e^{-\Omega(\alpha^2 m)} \) that
\[
\| A_f(r_1 - r_2) \|_2 \geq \sqrt{m}(\mu - \alpha)\| r_1 - r_2 \|_2.
\]
Taking a union bound over \( G(M) \times G(M) \), we have that when \( m = \Omega(\frac{k}{\alpha^2} \log \frac{Lr}{\delta}) \), with probability \( 1 - e^{-\Omega(\alpha^2 m)} \), for all \( r_1, r_2 \in G(M) \),
\[
\| A_f(r_1 - r_2) \|_2 \geq \sqrt{m}(\mu - \alpha)\| r_1 - r_2 \|_2.
\]
Therefore, we have that when \( m = \Omega(\frac{k}{\alpha^2} \log \frac{Lr}{\delta}) \), with probability \( 1 - e^{-\Omega(\alpha^2 m)} \),
\[
\frac{1}{\sqrt{m}} \| A_f(s_0 - t_0) \|_2 \geq (\mu - \alpha)\| s_0 - t_0 \|_2 \geq (\mu - \alpha)(\| x_1 - x_2 \|_2 - 2\delta),
\]
where we applied (142). In addition, using the results in [5, 35], for the standard Gaussian matrix \( A \), we have \( m = \Omega(\frac{k}{\alpha^2} \log \frac{Lr}{\delta}) \), with probability \( 1 - e^{-\Omega(\alpha^2 m)} \),
\[
\sum_{i=1}^{\ell} (\| A(s_i - s_{i-1}) \|_2 + \| A(t_i - t_{i-1}) \|_2) \leq O(\sqrt{m\alpha\delta}) = O(\sqrt{m\delta}).
\]
Moreover, by [60], Corollary 5.35, we have \( \| \frac{1}{\sqrt{m}} A \|_{2 \rightarrow 2} \leq 2 + \sqrt{\frac{\pi}{m}} \) with probability at least \( 1 - e^{-m/2} \). Hence, choosing \( \ell = \lfloor \log_2 n \rfloor \), we have with probability at least \( 1 - e^{-m/2} \) that
\[
\left\| \frac{1}{\sqrt{m}} A(t_\ell - x_2) \right\|_2 + \left\| \frac{1}{\sqrt{m}} A(s_\ell - x_1) \right\|_2 \leq 2 + \sqrt{\frac{\pi}{m}} \frac{\delta}{2^\ell} = O(\delta),
\]
where we used the fact that \( \|t_\ell - x_2\| \leq \frac{\delta}{2^\ell} \) and \( \|s_\ell - x_1\| \leq \frac{\delta}{2^\ell} \). Combining (144), (147), (148) and (149), we obtain the desired result.

### E Further Experiments: Comparing the Iterative Algorithm in CoPRAM Combined with Various Initialization Methods

In this section, we provide additional numerical results comparing our approach with other initialization methods, in that case that the initialization is combined with the subsequent iterative algorithm used in CoPRAM. In addition to PRI-SPCA-NT and the spectral initialization methods used in ThWF, SPARTA, and CoPRAM, we also compare with random initialization, in which the initial vector is \( \lambda \frac{r}{\| r \|_2} \) with \( r \sim \mathcal{N}(0, I_n) \). The worsened performance of random initialization reveals the importance of using spectral initialization methods for sparse phase retrieval.

We first consider the noiseless case and compare the relative error and empirical success rate of different methods. More specifically, after obtaining initial vectors from different initialization methods using the procedure discussed in Section 5, we run the iterative algorithm of CoPRAM for \( T = 100 \) iterations to further refine the estimated vectors and obtain \( x^{(T)} \). For our experiments, we found that 100 iterations are usually sufficient for convergence. A trial is declared to be successful if the relative error
\[
\min\{ \| x^{(T)} - x \|_2, \| x^{(T)} + x \|_2 \} \frac{1}{\| x \|_2}
\]
is less than 0.01. We fix \( n = 1000 \), and consider \( s = 10 \) or \( s = 20 \), while varying \( m \) in \( \{100, 150, 200, \ldots, 1000\} \). For each of the experiments in this section, we repeat 50 trials for 10 times, and calculate the standard deviation over these 10 times to produce the error bars.

We report the relative error in Figure 4 and we report the empirical success rate of different methods in Figure 5. From Figures 4 and 5, we observe that for most cases, our PRI-SPCA method leads to the smallest relative error and the largest empirical success rate.

Next, we consider the noisy case, and compare the relative error for using approximately the same time cost. The time cost is calculated as the sum of the running time of each initialization method and the running time of the subsequent iterative algorithm. As we have mentioned in Section 5, the time complexity of each iteration in GRQI is \( O(s^3 + m) \), while for the iterative algorithm of CoPRAM, the time complexity of each iteration is \( O(s^2 n \log n) \), which dominates the total time complexity.
In this experiment, we fix $n = 1000$, $m = 500$, $s = 20$, and consider $\sigma = 0.1$ and $0.2$. Since the running times of the initialization methods are typically less than 0.1 second, we vary the time cost (in seconds) $t$ in $\{0.1, 0.12, \ldots, 0.48, 0.5\}$. For each $t$ and each of the methods we consider, we find the number of the iteration whose time cost is closest to $t$, and record the corresponding relative error. Note that as mentioned in Section 5, for noisy measurements, we do not compare with ThWF because it corresponds to quadratic measurements. The results are reported in Figure 6 from which we observe that when combined with the iterative algorithm of CoPRAM and using approximately the same time cost, our PRI-SPCA method gives smallest relative error in most cases.
Figure 6: Relative error for different initialization methods combined with the iterative algorithm of CoPRAM in the noisy setting with $\sigma = 0.1$ (Left) and $\sigma = 0.2$ (Right), when using approximately the same time cost.