Universal quantum computation using only projective measurement, quantum memory, and preparation of the $|0\rangle$ state

Michael A. Nielsen*

Centre for Quantum Computer Technology, University of Queensland, Queensland 4072, Australia.
(October 25, 2018)

What resources are universal for quantum computation? In the standard model, a quantum computer consists of a sequence of unitary gates acting coherently on the qubits making up the computer. This paper shows that a very different model involving only projective measurements, quantum memory, and the ability to prepare the $|0\rangle$ state is also universal for quantum computation. In particular, no coherent unitary dynamics are involved in the computation.

Recall that in the standard quantum circuits model (see, for example, [1] for a review) a quantum computation consists of three elements: (a) an initialization stage in which the computer is prepared in an $n$-qubit computational basis state; (b) a processing stage in which a sequence of one- and two-qubit unitary gates is applied to the computer; and (c) a read-out stage in which the result of the computation is read out by measuring some subset of the qubits in the computational basis.

This paper demonstrates that a surprising alternate model is also universal for quantum computation. In this measurement model for quantum computation only the following three operations are allowed: (a) preparation of qubits in the $|0\rangle$ state; (b) storage of qubits (quantum memory); and (c) non-destructive) projective measurements on up to four qubits at a time, in an arbitrary basis. What is surprising about the universality of this model is that no coherent dynamical operations are allowed, contrary to the widespread folklore belief that such operations are crucial to universal quantum computation.

Substantial prior work has been done on the physical requirements for universal quantum computation. Barenco et al [2] showed that controlled-NOT and single-qubit unitary operations are sufficient to do universal quantum computation. Deutsch, Ekert and Barenco [3] and Lloyd [4] showed that almost any two-qubit quantum gate is universal. More recently, Dodd et al [5] have shown that any two-body Hamiltonian entangling qubits, together with local unitary operations, is universal for computation. (See also related work in [6].)

Thus, the conventional approach to universality has been to identify a set of coherent dynamical operations universal for quantum computation. Recently, however, it has been realized that quantum measurement is a powerful primitive element that can be performed during a computation. For example, Raussendorf and Briegel [6] have shown that by combining quantum measurements with the ability to prepare a special “cluster” state using a coherent Ising-type interaction, it is possible to do universal quantum computation. Knill, Laflamme and Milburn [7] have shown that single-photon detection and single-photon sources enable universal quantum computation in optics, using otherwise relatively easy coherent dynamical operations from linear optics.

In this paper I extend these results to show that no coherent dynamics at all are required to do universal quantum computation. The results of this paper build on a method developed by Nielsen and Chuang [8], who showed how to stochastically “teleport” a quantum gate from one location to another. Elegant generalizations of this method have been developed in [9–11]. The method has also been applied by Gottesman and Chuang [12] to develop fault-tolerant constructions for quantum gates, and further work [13] has been done simplifying such constructions. Note also that the optical quantum computer proposed by Knill, Laflamme and Milburn [12] is based on the gate teleportation idea.

In order to show that the measurement model can simulate the standard model of quantum computation we need only show that the measurement model can simulate the controlled-NOT gate, as well as any single-qubit unitary gate. We begin by explaining how the measurement model can simulate the action of an arbitrary single-qubit gate $U$ on a single-qubit state $|\psi\rangle$.

The first step is to use the measurement model to prepare one of the two-qubit states $|U_j\rangle$ defined by $|U_j\rangle \equiv (I \otimes U \sigma_j)(|00\rangle + |11\rangle)/\sqrt{2}$, where $\sigma_j$ are the four Pauli $\sigma$ matrices, $I$, $\sigma_x$, $\sigma_y$, and $\sigma_z$. Simple algebra shows that these states form an orthonormal basis for the state space of two qubits, so we can achieve this state preparation by first preparing the state $|00\rangle$, and then measuring in the orthonormal basis of states $|U_j\rangle$. It is important to note that the construction presented below works no matter which of the four states $|U_j\rangle$ is output from this measurement procedure.

Having prepared the state $|U_j\rangle$ offline, we now attempt to use this state and Bell-basis measurements to perform the operation $U$ on $|\psi\rangle$. The basic operation required to achieve this is the gate teleportation circuit shown in...
Following [13], to analyse the output from the third line of this circuit, it is convenient to note that $U\sigma_j$ acting on the third qubit commutes with the Bell measurement on the first two qubits, so Fig. 1 is equivalent to Fig. 3. The circuit in Fig. 3 is easy to understand; except for the final gate $U\sigma_j$ acting on the third line, it is just quantum teleportation [21], and thus the final state output from the third line is $U\sigma_j\sigma_m|\psi\rangle$, where $\sigma_m$ is one of the four Pauli matrices determined by the measurement result $m$ from the Bell measurement. Note that each of the four possible outcomes $U\sigma_j|\psi\rangle, U\sigma_j\sigma_x|\psi\rangle, U\sigma_j\sigma_y|\psi\rangle$, and $U\sigma_j\sigma_z|\psi\rangle$ occurs with equal probability 1/4.

Thus, with probability 1/4 the result $m$ is the same as $j$, and the gate teleportation succeeds, with the gate $U$ being applied to the qubit. With probability 3/4, however, $m \neq j$, and the incorrect operation $U\sigma_j\sigma_m$ is applied to the qubit. At first glance this appears to create a significant problem, since we can’t just discard the qubit and start again, as it may have been in an unknown state produced as part of a quantum computation. Fortunately, this problem can be avoided by teleporting the gate $U\sigma_m\sigma_jU^\dagger$. If this succeeds, which occurs with probability 1/4, then the net action is $U\sigma_m\sigma_jU^\dagger U\sigma_j\sigma_m = U$, as desired. If the second gate teleportation fails, then the net action is $U\sigma_m\sigma_jU^\dagger \sigma_{m'}\sigma_j' U\sigma_j\sigma_m$, where $|U_{j'}\rangle$ is the state that was prepared to do the second gate teleportation, and $m'$ is the measurement result from the second gate teleportation. This is still okay, as we can use a similar gate teleportation procedure again to attempt to obtain the correct dynamics $U$. In general, the procedure we use to simulate $U$ in the measurement model is as follows:

1. **Initialization**
   - $r := 1$; loop counter
   - $C_0 := I$; cumulative effect after 0 iterations

2. **Main Loop**
   - $A^r := UC_{r-1}^i$; the gate we attempt to teleport
   - Prepare $|A^r_{j_r}\rangle$
   - Teleport $A^r$, return result $m_r$
   - $C_r := A^r\sigma_j, \sigma_m, C_{r-1}$; cumulative effect
   - **Case:** $m_r = j_r$; success
   - **Case:** $m_r \neq j_r$; failure
   - $r := r + 1$; update loop counter
   - **Goto Main Loop**; try again

This procedure for simulating $U$ has an intrinsic error probability due to the possibility of failure in the gate teleportation procedure. If we demand that the procedure succeed with probability one, then the procedure may be repeated an arbitrarily large number of times, making our simulation inefficient. Fortunately, an efficient construction can be devised, based on the threshold theorem for quantum computation [21–26]. The idea is to use the measurement model to simulate a fault-tolerant circuit in the standard model. In such a fault-tolerant circuit any element in the circuit can fail with some small probability $\epsilon > 0$ (currently estimates put $\epsilon$ in the range $10^{-4}$ to $10^{-6}$), yet the circuit as a whole still succeeds with probability arbitrarily close to one. The single-qubit gates in the fault-tolerant circuit can thus be simulated in the measurement model, using at most $r$ iterations of the gate teleportation procedure, where $(3/4)^r \epsilon < \epsilon$. Thus the total number of operations required to simulate $U$ in the measurement model scales as $O(\log(1/\epsilon))$. If we assume a threshold of $10^{-5}$ then we require $r = 41$ iterations to achieve a failure probability less than $10^{-5}$.

Simulating two-qubit gates in the measurement model of quantum computation is similar to simulation of single-qubit gates. Let $U$ now denote an arbitrary two-qubit quantum gate. To simulate $U$ we first define $|U_{jk}\rangle$ to be the result of applying $U(\sigma_j \otimes \sigma_k)$ to the third and fourth qubits of $(|00\rangle + |11\rangle) \otimes (|00\rangle + |11\rangle)/2$, where the subscripts label which qubit is being referred to. Such a state can be prepared by first preparing the state $|00\rangle$ and then measuring in the orthonormal basis $|U_{jk}\rangle$. Gate teleportation is now achieved using Fig. 1, in a fashion analogous to the single-qubit case. Once again, the number of operations necessary to achieve a failure probability at most $\epsilon > 0$ is $O(\log(1/\epsilon))$ operations in the measurement model.
The simplicity of the model may also make it useful for the study of quantum computational complexity. Finally, by showing that coherent dynamical operations are not necessary for universal quantum computation, the model gives valuable insight into the fundamental question of what gives quantum computers their power.

**ACKNOWLEDGMENTS**

This paper was stimulated by a discussion with Damian Pope about what makes quantum computers powerful. Thanks to Michael Bremner, Ike Chuang, Jennifer Dodd, Alexei Gilchrist, Tobias Osborne, and Damian Pope for a thorough reading of the manuscript.

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