Diameter of the Stochastic Mean-Field Model of Distance

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We consider the complete graph $K_n$ on $n$ vertices with exponential mean $n$ edge lengths. Writing $C_{ij}$ for the weight of the smallest-weight path between vertices $i, j \in [n]$, Janson [18] showed that $\max_{i,j \in [n]} C_{ij} / \log n$ converges in probability to 3. We extend these results by showing that $\max_{i,j \in [n]} C_{ij} - 3 \log n$ converges in distribution to some limiting random variable that can be identified via a maximization procedure on a limiting infinite random structure. Interestingly, this limiting random variable has also appeared as the weak limit of the re-centred graph diameter of the barely supercritical Erdős–Rényi random graph in [22].

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1. Introduction

We consider the complete graph $\mathcal{K}_n$ on the vertex set $[n] := \{1, 2, \ldots, n\}$ and edge set $\mathcal{E}_n := \{\{i, j\} : i < j \in [n]\}$. Each edge $e \in \mathcal{E}_n$, is assigned an exponential mean $n$ edge length $E_e$, independent across edges. This implies that for any vertex $v$, the closest neighbour to this vertex is $O(1)$ distance away. Define the length of a path $\pi$ as

$$w(\pi) := \sum_{e \in \pi} E_e. \quad (1.1)$$

This assignment of random edge lengths makes $\mathcal{K}_n$ a (random) metric space often referred to as the stochastic mean-field model of distance (see Section 3). By continuity of the distribution of edge lengths, this metric space has unique geodesics. For any two vertices $i, j \in [n]$, let $\pi(i, j)$ denote the shortest path between these two vertices and write $C_{ij}$ for the length of this geodesic; precisely, writing $\mathcal{P}_{ij}$ for the collection of all paths from $i$ to $j$, then $C_{ij} = \min_{\pi \in \mathcal{P}_{ij}} w(\pi)$. The functional of interest in this paper is the diameter of the metric space:

$$\text{Diam}_n(\mathcal{K}_n) := \max_{i,j \in [n]} C_{ij}. \quad (1.2)$$
We first dive into the statement of the main result, postponing a full discussion to Section 3.

2. Results

The main aim of this paper is to prove that the diameter defined in (1.2) properly re-centred converges to a limiting random variable. We start by constructing this limiting random variable.

2.1. Construction of the limiting random variable

The limiting random variable arises as an optimization problem on an infinite randomly weighted graph $G_{\infty} = (V, E)$. The vertex set of this graph is the set of positive integers $\mathbb{Z}_+ = \{1, 2, \ldots \}$, while the edge set consists of all undirected edges $E = \{(i, j) : i, j \in \mathbb{Z}_+, i \neq j \}$. Let $P$ be a Poisson process on $\mathbb{R}$ with intensity measure having density

$$\lambda(y) = e^{-y}, \quad -\infty < y < \infty. \quad (2.1)$$

It is easy to check that $\max \{x : x \in P \} < \infty$ a.s. Thus we can order the points in $P$ as $Y_1 > Y_2 > \cdots$. We think of $Y_i$ as the vertex weight at $i \in \mathbb{Z}_+$. The edge weights are easier to describe. Let $(\Lambda_{st})_{s,t \in \mathbb{Z}_+, s < t}$ be a family of independent standard Gumbel random variables, namely $\Lambda_{st}$ has cumulative distribution function

$$F(x) = e^{-e^{-x}}, \quad -\infty < x < \infty. \quad (2.2)$$

The random variable $\Lambda_{s,t}$ gives the weight of an edge $\{s, t\} \in E$. Now consider the optimization problem

$$\Xi := \max_{s,t \in \mathbb{Z}_+, s < t} (Y_s + Y_t - \Lambda_{st}). \quad (2.3)$$

Although this is not obvious, we shall show that $\Xi < \infty$ a.s. The main result in this paper is as follows. We write $\overset{w}{\rightarrow}$ to denote convergence in distribution.

**Theorem 2.1 (diameter asymptotics).** For the diameter of the stochastic mean-field model of distance, as $n \to \infty$,

$$\max_{i,j \in [n]} C_{ij} - 3 \log n \overset{w}{\rightarrow} \Xi,$$

and

$$\mathbb{E}\left[ \max_{i,j \in [n]} C_{ij} \right] - 3 \log n \to \mathbb{E}[\Xi], \quad \text{Var}\left( \max_{i,j \in [n]} C_{ij} \right) \to \text{Var}(\Xi). \quad (2.4)$$

**Remark.** Theorem 2.1 solves [18, Problems 1 and 2]. Further, our proof shows that for any fixed $p \geq 1$, as $n \to \infty$

$$\mathbb{E}\left( \max_{i,j \in [n]} C_{ij} - 3 \log n \right)^p \to \mathbb{E}(\Xi^p).$$

2.2. Basic notation

Let us briefly describe the notation used in the rest of the paper. We write $\overset{P}{\rightarrow}$ to denote convergence in probability. For a sequence of random variables $(X_n)_{n \geq 1}$, we write $X_n = O_P(b_n)$
when $(|X_n|/b_n)_{n \geq 1}$ is a tight sequence of random variables as $n \to \infty$, and $X_n = o_p(b_n)$ when $|X_n|/b_n \overset{p}{\to} 0$ as $n \to \infty$. Further we write $X_n = \Theta_P(b_n)$ if $(X_n/b_n)_{n \geq 1}$ and $(b_n/X_n)_{n \geq 1}$ are both a tight sequence of random variables. For a non-negative function $n \to g(n)$, we write $f(n) = O(g(n))$ when $|f(n)|/g(n)$ is uniformly bounded, and $f(n) = o(g(n))$ when $\lim_{n\to \infty} f(n)/g(n) = 0$. Furthermore, we write $f(n) = \Theta(g(n))$ if $f(n) = O(g(n))$ and $g(n) = O(f(n))$. Finally, we write that a sequence of events $(A_n)_{n \geq 1}$ occurs with high probability (w.h.p.) when $P(A_n) \to 1$. We use $Y \sim \exp(\lambda)$ for a random variable with exponential distribution with rate $\lambda$. At several places we will use the fact that for $Y \sim \exp(1)$ then $\log(1/Y)$ has a standard Gumbel distribution with distribution as in (2.2). To ease notation we will occasionally use symbols representing a real number when an integer is required (e.g. $\sum_{j=1}^{\sqrt{n}}$) and suppress the precise $\lfloor \cdot \rfloor$ notation.

3. Background and related results

We now discuss our results and place them in the context of results in the literature.

3.1. Stochastic mean-field model of distance

The stochastic mean-field model of distance has arisen in a number of different contexts in understanding the structure of combinatorial optimization problems in the presence of random data, ranging from shortest path problems [18] to random assignment problems [2, 3], minimal spanning trees [15, 17] and travelling salesman problems [25]; see [6] for a comprehensive survey and related literature. The closest work to this study is the paper by Janson [18]. Recall that $C_{ij}$ denotes the length of the geodesic between two vertices $i, j \in [n]$; by symmetry this has the same distribution for any two vertices in $i, j$. For any vertex $i \in [n]$, write $\text{Flood}[i] := \max_{j \in [n]} C_{ij}$ for the maximum time started at $i$ to reach all vertices in $\mathcal{K}_n$ (often called the flooding time). Then Janson proved that, as $n \to \infty$,

$$\frac{C_{ij}}{\log n} \overset{p}{\to} 1, \quad \frac{\text{Flood}[i]}{\log n} \overset{p}{\to} 2, \quad \frac{\text{Diam}_{\mathcal{K}_n}}{\log n} \overset{p}{\to} 3,$$  \hspace{1cm} (3.1)

and further

$$C_{ij} - \log n \overset{w}{\to} \Lambda_1 + \Lambda_2 - \Lambda_{12},$$ \hspace{1cm} (3.2)

while

$$\text{Flood}[i] - 2\log n \overset{w}{\to} \Lambda_1 + \Lambda_2.$$ \hspace{1cm} (3.3)

Here $\Lambda_1, \Lambda_2, \Lambda_{12}$ are all independent standard Gumbel random variables as in (2.2). Problems 1 and 2 in [18] then ask whether a result similar to (3.2) and (3.3) holds for the diameter $\text{Diam}_{\mathcal{K}_n}$ (with (3.1) obviously re-centred by $3\log n$).

The main aim of this paper is to answer this question in the affirmative. We discuss more results about the distribution of $\Xi$ in Section 4.8. In the context of (2.4), for $C_{ij}$ and $\text{Flood}[i]$, Janson also shows convergence of the expectation and variance with explicit limit constants. We have been unable to derive explicit values for the limit constants $E(\Xi)$ and $\text{Var}(\Xi)$. 


3.2. Hopcount and extrema

This paper looks at the length of optimal paths (measured in terms of the edge weights). One could also look at the hopcount or the number of edges $|\pi(i, j)|$ on the optimal path as well as the largest hopcount $D^* = \max_{i,j \in [n]} |\pi(i, j)|$. The entire shortest path tree from a vertex $i$ has the same distribution as a random recursive tree on $n$ vertices (see [24] for a survey). Janson used this in [18] to show that $|\pi(i, j)| - \log n \sqrt{\log n} \xrightarrow{w} Z$, where $Z$ has a standard normal distribution. The maximal hopcount $\mathcal{H}_n(i) = \max_{j \in [n]} |\pi(i, j)|$ from a vertex $i$ has the same distribution as the height of a random recursive tree, which by [12] or [21] satisfies the asymptotics $\mathcal{H}_n(i)/\log n \xrightarrow{p} e$ as $n \to \infty$.

The first-order asymptotics for the maximum hopcount $D^*$ were recently proved in [1], showing that $D^*/\log n \xrightarrow{p} \alpha^*$, where $\alpha^* \approx 3.5911$ is the unique solution of the equation $x \log x - x = 1$.

3.3. First passage percolation on random graphs

The last few years have seen progress in the understanding of optimal paths in the presence of edge disorder (usually assumed to have exponential distribution) in the context of various random graph models (see e.g. [7, 8, 11] and the references therein). In particular, Proposition 4.4 below with a sketch of proof has appeared in [4, 5, 10].

In the context of our main result, [7] studied the weighted diameter for the random $r$-regular graphs $G_{n,r}$ with exponential edge weights, and proved first-order asymptotics. We conjecture that one can adapt the main techniques in this paper to show the second-order asymptotics for $r \geq 3$, that is,

$$Diam_w(G_{n,r}) - \left(\frac{1}{r-2} + \frac{2}{r}\right) \log n \xrightarrow{w} \Xi_r,$$

for a limit random variable $\Xi_r$ that satisfies that, as $r \to \infty$,

$$r\Xi_r \xrightarrow{w} \Xi.$$  \hfill (3.5)

3.4. Diameter of the barely supercritical Erdős–Rényi random graph

Consider the barely supercritical Erdős–Rényi random graph $G_n(n, (1+\varepsilon)/n)$ where $\varepsilon = \varepsilon_n \to 0$ but $\varepsilon n^3 \to \infty$. It turns out that the random variable $\Xi$ in Theorem 2.1 is closely related to the random variable describing second-order fluctuations for the graph diameter $Diam_w(G_{n,r}(n, (1+\varepsilon)/n))$. Here we use $Diam_w(\cdot)$ for the graph diameter of a graph, namely the largest graph distance between any two vertices in the same component. We now describe this result. Consider the minor modification of the optimization problem defining $\Xi$ in Section 2 where the Poisson
process $\mathcal{P}$ generating the vertex weights has intensity measure with density

$$\lambda(y) = ye^{-y}, \quad -\infty < y < \infty.$$  

As before, the edge weights $\Lambda_{uv}$ are independent standard Gumbel random variables. Let $\Xi_\gamma$ denote the random variable corresponding to the optimization problem in (2.3). Let $\lambda = 1 + \epsilon$ and let $\lambda_* < 1$ be the unique value satisfying $\lambda_* e^{-\lambda_*} = \lambda e^{-\lambda}$. After an initial analysis in [13, 14], Riordan and Wormald in [22, Theorem 5.1] showed that there exists a constant $\gamma > 0$ such that

$$\text{Diam}_p(G_n(n(1+\epsilon)/n)) - \frac{\log \epsilon^3 n}{\log \lambda} - \frac{2 \log \epsilon^3 n}{\log 1/\lambda_*} \xrightarrow{\text{w}} \Xi_\gamma.$$  

We believe that the Poisson cloning technique in [13, 14] coupled with the techniques in this paper may yield an alternative proof of this result, but we defer this to future work.

4. Proofs

We start with the basic ideas behind the main result. We then describe the organization of the rest of the section, which deals with converting this intuitive picture into a proper proof.

4.1. Proof idea

We write $\mathcal{S}_n = (\mathcal{K}_n, (E_e)_{e \in E_n})$ for the (random) metric space where $(E_e)_{e \in E_n}$ are i.i.d. mean $n$ exponential random variables. Recall from Janson’s result, as described in (3.2), that the distance $C_{ij}$ between typical vertices $i, j \in [n]$ scales like $\log n + O_{\nu}(1)$. Intuitively, the extra $2\log n$ in the diameter arises due to the following reason. Consider ranking the vertices according to the distance to their closest neighbour. More precisely, for each vertex $i \in [n]$, write $X_{i} = \min_{j \neq i} (\epsilon^{-1} E_{ij})$, the distance to the closest vertex to $i$. Arrange these as $X_{(V_1)} > X_{(V_2)} > \cdots > X_{(V_n)}$.

We shall show that:

(a) the point process $\mathcal{P}_n = (X_{(V_i)} - \log n)_{i \geq 1}$ converges to the Poisson point process $\mathcal{P}$ in Section 2 with intensity measure given by (2.1);

(b) the diameter of $\mathcal{K}_n$ corresponds to the shortest path between a pair of these ‘slow’ vertices $(V_s, V_l)$;

(c) further, after reaching the closest vertex, the remaining path behaves like a typical optimum path in the original graph $\mathcal{K}_n$ equipped with exponential mean $n$ edge lengths, but now between two disjoint pairs of vertices.

More precisely, part (c) entails that $C_{V_s V_l} \approx X_{(V_1)} + X_{(V_2)} + d_w(A,B)$, where $A = \{a,b\}$, $B = \{c,d\}$ with $a, b, c, d$ four distinct vertices in $[n]$ and $d_w(A,B)$ is a random variable independent of $X_{(V_1)}, X_{(V_2)}$ having the same distribution as the distance between the sets $A, B$ in the original metric space $\mathcal{S}_n$. The first two terms correspond to the time to get out of these ‘slow’ vertices, which scale like $\log n + O_{\nu}(1)$ by (a) while $d_w(A,B)$ scales like $\log n + O_{\nu}(1)$, thus implying that the diameter scales like $3\log n + O_{\nu}(1)$. By investigating the fluctuations of $X_{(V_1)}, X_{(V_2)}$ and $d_w(A,B)$, we can also identify the fluctuations of $\max_{i,j \in [n]} C_{ij}$.

**Organisation of the proof.** We start in Section 4.2 by describing the distribution of the shortest path between two disjoint sets of vertices. Section 4.3 proves a weaker version of the Poisson point process limit described in (a) above. Section 4.4 describes the limiting joint distribution of
the (properly re-centred) weights of optimal paths between multiple source destination pairs in \( \mathcal{S}_n := (\mathcal{K}_n, \{ E_e : e \in E_n \}) \). Section 4.5 uses the results in Section 4.3 and 4.4 to study asymptotics for the joint distribution of distances between the slow vertices \((V_s)_{s \in [n]}\). Section 4.6 shows that the diameter of \( \mathcal{K}_n \) corresponds to the optimal path between one of the ‘first few’ slow vertices. The last three sections use these ingredients to show distributional convergence as well as the convergence of the moments of \( \text{Diam}_w(\mathcal{K}_n) - 3\log n \) to the limiting random object, thus completing the proof of the main result.

4.2. Explicit distributions for distances between sets of vertices

In this section we explain the proof by Janson of (3.2). We also extend that analysis to the smallest-weight path between disjoint sets of vertices. We remind the reader that the standing assumption henceforth is that each edge has exponential mean \( n \) distribution. We start with the following lemma.

**Lemma 4.1 (distances between sets of vertices).** Consider two disjoint non-empty sets \( A, B \subseteq [n] \). Then,

\[
d_w(A, B) \stackrel{d}{=} \sum_{k=|A|}^{N+|A|-1} \frac{E_k}{k(n-k)},
\]

where

(i) \( (E_k)_{k \geq 1} \) are i.i.d. mean \( n \) exponential random variables;

(ii) \( N \) is independent of the sequence \( (E_k)_{k \geq 1} \) with the same distribution as the number of draws required to select the first black ball in an urn containing \( |B| \) black balls and \( n - |A| - |B| \) white balls, where one is drawing balls without replacement from the urn.

**Proof.** We start by exploring the neighbourhood of the set \( A \) in a similar way as in [18]. Recall that each edge has an exponential mean \( n \) edge length. After having found the \( \ell \)th minimal edge and with \( k = (|A| + \ell) \), there are \( k(n-k) \) edges incident to the found vertices. The minimal edge weight thus has an exponential distribution with mean \( n/k(n-k) \). This process stops the first time we find a vertex in \( B \). Since every new vertex added to the cluster of reached vertices is chosen uniformly amongst the set of present unreached vertices, the distribution of the number of steps required to reach a vertex in \( B \) has the distribution \( N \) asserted in the lemma, independently of the inter-arrival times of new vertices found. Thus the time it takes to find the first element in \( B \) is

\[
\sum_{\ell=0}^{N-1} \frac{E_k}{(\ell + |A|)(n-\ell - |A|)}.
\]

Defining \( k = \ell + |A| \) proves the claim. \( \square \)

Now we specialize to a particular case of the above lemma. Fix a vertex, say vertex \( v = 1 \), and another set \( B \subseteq [n] \setminus \{1\} \). For much of what follows, we will be concerned with the optimal path between such a vertex and a set of size \( |B| = \Theta(\sqrt{n}) \). This is an appropriate time to think about two different but equivalent ways to find such an optimal path.
**Process 1.** The first way to find the optimal path is the exploration process described in the previous lemma where we start at vertex $v = 1$ and keep adding the closest vertex to the cluster until we hit a vertex in $B$. Write $M_B$ for the number of vertices other than $B$ that are found in this exploration. Note that the proof of the previous lemma implies that
\[
(dw(\{1\}, B), M_B) \overset{d}{=} \left( \sum_{k=1}^{N_B} E_k / k(n-k), N_B \right),
\]
where $N_B$ is independent of the sequence $(E_k)_{k \geq 1}$ and has the same distribution as the number of balls required to get the first black ball when drawing balls without replacement from an urn containing $|B|$ black balls and $n - 1 - |B|$ white balls.

**Process 2.** The second way to find the optimal path is the following. We think of water starting at source vertex $v = 1$ at time $t = 0$ percolating through the network at rate one using the edge lengths. Write $SWG_t^{(1)}$ (an acronym for smallest-weight graph) for the set of vertices reached by time $t$ starting from vertex 1. More precisely,
\[
SWG_t^{(1)} := \{ u \in [n] : dw(1, u) \leq t \}.
\]
By convention, vertex $v = 1$ is in $SWG_t^{(1)}$ for all $t \geq 0$. Now note that the size process $(|SWG_t^{(1)}|)_{t \geq 0}$ is a pure-birth Markov process (with respect to the filtration $(\mathcal{F}_t)_{t \geq 0} = (\sigma(SWG_t))_{t \geq 0}$) with rate of birth given by $n/k(n-k)$ when the size $|SWG_t^{(1)}| = k$. Each new vertex added to this cluster is chosen uniformly amongst all available unreached vertices at that time, that is, the vertices $[n] \setminus SWG_t^{(1)}$. Finally, the distance $dw(\{1\}, B)$ can be recovered as
\[
dw(\{1\}, B) := \inf \{ t \geq 0 : SWG_t^{(1)} \cap B \neq \emptyset \}.
\]

In this section, we use Process 1 to prove the following initial result. We use Process 2 in Section 4.4 below.

**Lemma 4.2 (distances between vertex and set of size $b\sqrt{n}$).** Let $B \subseteq [n]$ with $|B| = b\sqrt{n}$. Then as $n \to \infty$,
\[
\left( dw(\{1\}, B) - \frac{1}{2} \log n, M_B / \sqrt{n} \right) \overset{w}{\to} (\Lambda + \log (\hat{E} / b), \hat{E} / b),
\]
where $\hat{E}$ is exponential with parameter 1, $\Lambda$ is Gumbel, and $\hat{E}$ and $\Lambda$ are independent.

**Proof.** The above is equivalent to showing
\[
(dw(\{1\}, B) - \log M_B, M_B / \sqrt{n}) \overset{w}{\to} (\Lambda, \hat{E} / b),
\]
with $\Lambda, \hat{E}$ independent standard Gumbel and $\exp(1)$ respectively. Fix constants $0 < \alpha < \beta$ and $y \in \mathbb{R}$. Define the event
\[
A_n(y, \alpha, \beta) := \{ dw(\{1\}, B) - \log M_B \leq y \} \cap \{ \alpha \leq M_B / \sqrt{n} \leq \beta \}.
\]
Let \((E'_k)_{k \geq 1}\) be independent sequence of *mean one* exponential random variables. Equation (4.3) implies

\[
P(A_n(y, \alpha, \beta)) = \sum_{j=\alpha \sqrt{n}}^{\beta \sqrt{n}} P\left( \sum_{k=1}^{j} \frac{nE'_k}{k(n-k)} - \log j \leq y \right) P(N_B = j). \tag{4.7}
\]

Noting that \(\sum_{k=1}^{j} \frac{1}{j} \approx \log j + \gamma\) as \(j \to \infty\), where \(\gamma\) is Euler’s constant, gives

\[
\sum_{k=1}^{j} \frac{nE'_k}{k(n-k)} - \log j \approx \sum_{k=1}^{j} \frac{E'_k - 1}{k} + \gamma + R_n, \tag{4.8}
\]

where the error term \(R_n\) can be bounded independently of \(j\) by

\[
|R_n| \leq \sum_{k=1}^{j} \frac{E'_k}{n-k} \xrightarrow{p} 0, \tag{4.9}
\]

as \(n \to \infty\). Thus, uniformly for \(j \in [\alpha \sqrt{n}, \beta \sqrt{n}]\),

\[
P\left( \sum_{k=1}^{j} \frac{nE'_k}{k(n-k)} - \log j \leq y \right) \rightarrow P\left( \sum_{k=1}^{\infty} \frac{E'_k - 1}{k} + \gamma \leq y \right). \tag{4.10}
\]

By (4.7), to complete the proof it is enough to show that

\[
P(\alpha \leq N_B/\sqrt{n} \leq \beta) \rightarrow P(\alpha \leq \hat{E}/b \leq \beta). \tag{4.11}
\]

This follows easily, since for any \(x > 0\)

\[
P(N_B > x\sqrt{n}) = \prod_{k=1}^{x\sqrt{n}} \left( 1 - \frac{b\sqrt{n}}{n-1-k} \right) \sim e^{-bx},
\]

as \(n \to \infty\). \(\square\)

4.3. Poisson limit for the number of vertices with large minimal edge weights.

The aim of this section is to understand the distribution of edges emanating from the slow vertices, namely the set of vertices for which the closest vertex is at distance \(\approx \log n\). For vertex \(i \in [n]\), let \(X_{(i)} = \min_{j \in [n]} E_{ij}\) denote the minimal edge weight emanating from a given vertex \(i \in [n]\). Fix \(\alpha \in \mathbb{R}\) and let \(N_n(\alpha) = \#\{i \in [n] : X_{(i)} \geq \log n - \alpha\}\) denote the number of vertices with minimal outgoing edge weight at least \(\log n - \alpha\). We prove the following Poisson limit for \(N_n(\alpha)\).

**Proposition 4.3 (number of vertices with large minimal edge weight).** As \(n \to \infty\),

\[
N_n(\alpha) \xrightarrow{w} N(\alpha), \tag{4.12}
\]

where \(N(\alpha)\) is a Poisson random variable with mean \(e^\alpha\). More precisely, for \(n\) large enough,

\[
d_{TV}(N_n(\alpha), N(\alpha)) \leq \frac{2(1 + \varepsilon_n)e^{2\alpha} \log n}{n}, \tag{4.12}
\]
where \( d_{\text{tv}} \) denotes the total variation distance and
\[
\varepsilon_n = \exp \left( \frac{\log n - \alpha}{n} \right) - 1.
\]

**Proof.** We use the Stein–Chen method for Poisson approximation. Write
\[
N_n(\alpha) = \sum_{i \in [n]} Z_i, \quad Z_i = 1\{X_{(i)} \geq \log n - \alpha\}.
\]

For fixed \( i \in [n] \), note that \( X_{(i)} \) has an exponential distribution with mean \( n/(n - 1) \). Writing \( p_n = P(Z_i = 1) \) so that \( \lambda := E(N_n(\alpha)) = np_n \), it is easy to check that
\[
E(N_n(\alpha)) = (1 + \varepsilon_n)e^{\alpha}.
\] (4.13)

Thus, \( \lambda \to e^\alpha \) as \( n \to \infty \). For each fixed \( i \in [n] \), suppose we can couple \( N_n(\alpha) \) with a random variable \( W'_i \) such that the marginal distribution of \( W'_i \) is
\[
W'_i + 1 \overset{d}{=} N_n(\alpha) \mid \{Z_i = 1\},
\] (4.14)

that is, \( W'_i + 1 \) has the same distribution as \( N_n(\alpha) \) conditionally on \( \{Z_i = 1\} \). Then Stein–Chen theory [9] implies that in total variation distance
\[
d_{\text{tv}}(\mathcal{L}(N_n(\alpha)), \text{Poi}(\lambda)) \leq (1 + \lambda^{-1}) \sum_{i \in [n]} E(\varepsilon)^i E(|N_n(\alpha) - W'_i|)
\] (4.15)

Let us describe \( W'_i \), the same construction switching indices works for any \( i \). Let
\[
\mathcal{S}_n := \mathcal{K}_n, (E_e)_{e \in E_n}
\]
be the original edge lengths and let \( N_n(\alpha) \) be defined as above for the random metric space \( \mathcal{S}_n \). Let us construct the edge lengths of \( \mathcal{K}_n \) conditional on the event \( \{Z_1 = 1\} \) so that \( X_{(i)} - \log n \geq -\alpha \). We shall write \( \mathcal{S}'_n := \{(\mathcal{K}_n, (E'_e)_{e \in E_n})\} \) for \( \mathcal{S}_n \) conditioned on this event. Note that this event only affects edges incident to vertex 1 and further, by the lack of memory property of the exponential distribution, every such edge incident to vertex 1 has distribution \( \log n - \alpha + E \) where \( E \) is an exponential mean \( n \) random variable, independently across edges. Thus, we can construct the edge lengths on \( \mathcal{S}'_n \) using the edge lengths \( E_e \) in \( \mathcal{S}_n \) with the following description:

(a) for each edge \( e = \{1, i\} \) incident to vertex \( i \), set \( E'_e = \log n - \alpha + E_e \),

(b) for any edge not incident to vertex 1, set \( E'_e = E_e \).

Define \( X'_{(i)} \) analogously to \( X_{(i)} \) as the minimal edge length incident to vertex \( i \) but in \( \mathcal{S}'_n \). Finally, define
\[
Z'_i := 1\{X'_{(i)} > \log n - \alpha\}, \quad W'_i := \sum_{i \neq 1} 1\{X'_{(i)} \geq \log n - \alpha\}.
\]

Then \( W'_i \) by construction has the required distribution in (4.14). Note that
\[
|N_n(\alpha) - W'_i| \leq 1\{X_{(i)} > \log n - \alpha\} + \sum_{i \neq 1} |Z_i - Z'_i|.
\]

Taking expectations, by symmetry,
\[
E(|N_n(\alpha) - W'_i|) \leq p_n + (n - 1) E|Z_2 - Z'_2|.
\] (4.16)
Now
\[
E |Z_2 - Z'_2| = P(Z_2 = 1, Z'_2 = 0) + P(Z_2 = 0, Z'_2 = 1).
\]
Since the edge lengths in \( \mathcal{S}'_n \) are at least as large as the edge lengths in \( \mathcal{S}_n \), we have \{\( Z_2 = 1, Z'_2 = 0 \)\} = \( \emptyset \). For the second term,
\[
\{Z_2 = 0, Z'_2 = 1\} \equiv \{E_{2,1} < \log n - \alpha, \min_j E_{2,j} \geq \log n - \alpha\}.
\]
Since \( E_{i,j} \) are exponential mean \( n \), we immediately get
\[
P(Z_2 = 0, Z'_2 = 1) \leq (1 + o(1))e^{\alpha \log n}.\]
Using this in (4.16), the total variation bound (4.15) completes the proof.

4.4. Joint convergence of distances between multiple vertices
The aim of this section is to understand the re-centred asymptotic joint distribution of the minimal weight between multiple vertices. To prove this, it turns out that Process 2 using the smallest-weight graph \( \text{SWG}_v^{(t)} \) from vertices \( v \in [n] \) is more useful than Process 1. Versions of Proposition 4.4 below have appeared in [4, 5, 10]. We give a new proof, both for completeness and since we need a variant of this argument below.

Fix \( m \geq 2 \). Let \( (\Lambda_\alpha)_{\alpha \in [m]} \) and \( (\Lambda_{\alpha\beta})_{\alpha, \beta \in [m], \alpha < \beta} \) be independent standard Gumbel random variables. In the following proposition, we identify the limiting distribution of \( (d_w(\alpha, \beta) - \log n)_{\alpha, \beta \in [m], \alpha < \beta} \), an extension of the result given in (3.2) proved by Janson [18] for \( m = 2 \).

**Proposition 4.4 (joint distances between many vertices).** As \( n \to \infty \),
\[
(d_w(\alpha, \beta) - \log n)_{\alpha, \beta \in [m], \alpha < \beta} \xrightarrow{w} (\Lambda_\alpha + \Lambda_\beta - \Lambda_{\alpha\beta})_{\alpha, \beta \in [m], \alpha < \beta}.
\]

**Proof.** Fix \( m \geq 2 \). Write
\[
D(m) := (\Lambda_\alpha + \Lambda_\beta - \Lambda_{\alpha\beta})_{\alpha, \beta \in [m], \alpha < \beta},
\]
for the limiting array. The idea of the proof is as follows. We start by sequentially growing the smallest-weight graphs, SWGs, from the \( m \) vertices until they meet. This gives us a sequence of collision times \( (T_{\alpha\beta})_{\alpha < \beta \in [m]} \). An appropriately chosen linear transformation of these collision times stochastically dominates the array of the lengths of shortest paths. We show that this linear transformation of the collision times converges to the array \( D(m) \). A simple limiting argument using the convergence of the marginal distribution of two point distances implies that the joint distribution of the distances themselves converge to \( D(m) \) and this completes the proof.

Let us now begin the proof. We write \( \mathcal{S}_n \) for the random metric space \( (K_n, (E_e)_{e \in E_n}) \), where once again we remind the reader that \( E_e \) are i.i.d. exponential random variables with mean \( n \). Now start the smallest weight cluster \( \text{SWG}^{(t)}_v \) from vertex \( \alpha = 1 \). Write
\[
T_1 = \inf \{t : |\text{SWG}^{(t)}_v| = \sqrt{n}\}
\]
for the time for $\text{SWG}_{t}^{(1)}$ to grow to size $\sqrt{n}$. Now

$$ T_1 \overset{d}{=} \sum_{k=1}^{\sqrt{n}} nE_k/[k(n-k)], $$

where $\{E_k\}_{k \geq 1}$ is a sequence of independent rate one exponential random variables. This implies (see (4.8) and (4.10)) that

$$ T_1 - \frac{1}{2} \log n \overset{w}{\to} \log(1/\hat{E}_1), $$

(4.20)

where $\hat{E}_1$ is exponential with mean 1. For every vertex $v \in \text{SWG}_{i}^{(1)}$, write $B^{(1)}(v) := d_w(1,v)$ for the time when the flow from vertex 1 reaches $v$. We now work conditionally on the flow cluster $\text{SWG}_{i}^{(1)}$. By construction, as $n \to \infty$,

$$ \mathbb{P}(2 \notin \text{SWG}_{i}^{(1)}) = 1 - \sqrt{n} \to 1. $$

(4.21)

Further, by the memoryless property of the exponential distribution, conditionally on $\text{SWG}_{i}^{(1)}$, for every boundary edge $e = \{u,v\}$ with $u \in \text{SWG}_{i}^{(1)}$ and $v \notin \text{SWG}_{i}^{(1)}$, the remaining edge length $E_e - (T_1 - B^{(1)}(u))$ has an exponential distribution with mean $n$, and all these remaining edge lengths are independent.

Freeze the cluster $\text{SWG}_{i}^{(1)}$. Start a flow from vertex 2 as the source and write $\text{SWG}_{i}^{(2)}$ for the smallest-weight graph. Write

$$ T_{12} := \inf\{t : \text{SWG}_{i}^{(2)} \cap \text{SWG}_{i}^{(1)} \neq \emptyset\}, $$

(4.22)

so that $T_{12}$ is the first time that a vertex in the flow cluster from vertex $\alpha = 1$ at time $T_1$ is hit by the flow cluster from 2. Conditionally on $\text{SWG}_{i}^{(1)}$, on the event $\{2 \notin \text{SWG}_{i}^{(1)}\}$, we have that

(a) the length of the smallest-weight path between 1 and 2 is given by $d_w(1,2) = T_1 + T_{12}$;
(b) the random variable $T_{12}$ has the same distribution as $d_w(\{1\}, B)$ in the random (unconditional) metric space $S_n$, where $B$ is a fixed set of size $\sqrt{n}$.

By Lemma 4.2 with $b = 1$ we immediately get

$$ \left( T_{12} - \frac{1}{2} \log n, |\text{SWG}_{i}^{(2)}_{T_{12}}|/\sqrt{n} \right) \overset{w}{\to} (\log(1/\hat{E}_2) + \log(\hat{E}_{12}), \hat{E}_{12}), $$

(4.23)

where $\hat{E}_2$ and $\hat{E}_{12}$ are independent of each other and independent of $\hat{E}_1$, which arises as the limit in (4.20). Combining (4.20) and (4.23) we get

$$ (d_w(1,2) - \log n, |\text{SWG}_{i}^{(2)}_{T_{12}}|/\sqrt{n}) = \left( T_{12} - \frac{1}{2} \log n + T_1 - \frac{1}{2} \log n, N/\sqrt{n} \right) \overset{w}{\to} (\log(1/\hat{E}_1) + \log(1/\hat{E}_2) + \log(\hat{E}_{12}), \hat{E}_{12}). $$

(4.24)

This proves the claim for $m = 2$. We next extend the computation to $m = 3$.

For ease of notation, write $B = \sqrt{n} = |\text{SWG}_{i}^{(1)}|$ and $R = |\text{SWG}_{i}^{(2)}_{T_{12}}|$, here $B$ and $R$ will be mnemonics for ‘black’ and ‘red’ respectively. We now work conditionally on $A := \text{SWG}_{i}^{(1)} \cup \text{SWG}_{i}^{(2)}_{T_{12}}$. Since $|A| = \Theta_P(\sqrt{n})$,

$$ \mathbb{P}(3 \notin \text{SWG}_{i}^{(1)} \cup \text{SWG}_{i}^{(2)}_{T_{12}}) \to 1 \quad \text{as } n \to \infty. $$

(4.25)
Freeze the above two flow clusters. Start a flow from vertex $\beta = 3$ and consider the smallest-weight graph $SWG^{(3)}_3$ emanating from vertex 3. We need to modify this process after the first time it finds a vertex in $\mathcal{A} = SWG^{(3)}_{t_1} \cup SWG^{(2)}_{t_{12}}$, namely after time

$$T^*_3 = \inf\{ t : SWG^{(3)}_3 \cap \mathcal{A} \neq \emptyset \}.$$  

Suppose this happens due to $SWG^{(3)}_3$ finding a vertex in $SWG^{(1)}_{t_1}$. Remove all vertices in $SWG^{(1)}_{t_1}$ and all adjacent edges from $\mathcal{K}_n$ and then continue until the process finds a vertex in $SWG^{(3)}_{t_{12}}$. Similarly, if this happens due to a vertex in $SWG^{(2)}_{t_{12}}$ being found, then remove all vertices in $SWG^{(2)}_{t_{12}}$ and continue. Although this is not quite the smallest-weight graph emanating from vertex 3, to minimize notational overhead, we shall continue to denote this modified process by the same $SWG$. Remove all vertices in $\mathcal{A} \subset \{3\}$, to minimize notational overhead, we shall continue to denote this modified process by the same $SWG^{(3)}_3$. Define the stopping times

$$T_{13} = \inf\{ t \geq 0 : SWG^{(3)}_3 \cap SWG^{(1)}_{t_1} \neq \emptyset \},$$

and

$$T_{23} = \inf\{ t \geq 0 : SWG^{(3)}_3 \cap SWG^{(2)}_{t_{12}} \neq \emptyset \}.$$  

Similarly, define the sizes of the cluster $SWG^{(3)}_3$ at these stopping times as

$$C_n^{(13)} = |SWG^{(1)}_{t_1}|, \quad C_n^{(23)} = |SWG^{(2)}_{t_{12}}|.$$  

(4.26)

Similar to the urn description in (4.3), it is easy to check that conditionally on $\mathcal{A}$ and on the event $\{3 \notin \mathcal{A}\}$, the distribution of the random variables $(T_{13}, T_{23}, C_n^{(13)}, C_n^{(23)})$ can be constructed as follows. Consider an urn with $n$ balls out of which $B = |SWG^{(1)}_{t_1}|$ black balls, $R = |SWG^{(2)}_{t_{12}}|$ red balls and the remaining $n - B - R$ white balls. Also let $(E_k)_{k \geq 1}$ be an independent sequence of mean $n$ exponential random variables. Start drawing balls at random without replacement till the first time $\mathcal{N}_1$, we get either a black or a red ball.

(a) Suppose the first ball amongst the black or red balls is a black ball. Remove all black balls so that there are now $(n - \mathcal{N}_1 - B)$ balls in the urn. Continue drawing balls without replacement till we get a red ball. Let $\mathcal{N}_2 > \mathcal{N}_1$ be the time for the first pick of a red ball. Let $C_n^{(13)} = \mathcal{N}_1$, $C_n^{(23)} = \mathcal{N}_2$. Finally, let

$$T_{13} := \sum_{k=1}^{\mathcal{N}_1} \frac{E_k}{k(n-k)}, \quad T_{23} := T_{13} + \sum_{k=\mathcal{N}_1+1}^{\mathcal{N}_2} \frac{E_k}{k(n-k-B)},$$  

(4.27)

where, as before, $(E_k)_{k \geq 1}$ is an independent sequence of exponential random variables with mean $n$.

(b) Suppose the first ball amongst black and red balls to be picked is a red ball. Then, in the above formulae, simply interchange the roles of 1 and 2 and $B$ and $R$.

Using (4.23) and arguing exactly as in the proof of Lemma 4.2, we see that

$$\left( \frac{C_n^{(13)}}{\sqrt{n}}, \frac{C_n^{(23)}}{\sqrt{n}}, T_{13} - \frac{1}{2} \log n, T_{23} - \frac{1}{2} \log n \right) \rightarrow^w (\hat{E}_{13}, \hat{E}_{23}, \hat{E}_{12}, \log(1/\hat{E}_3) + \log(\hat{E}_{13}), \log(1/\hat{E}_3) + \log(\hat{E}_{23}/\hat{E}_{12})).$$  

(4.28)

Here $\hat{E}_3, \hat{E}_{13}, \hat{E}_{23}$ are independent of $\hat{E}_1, \hat{E}_2, \hat{E}_{12}$ and i.i.d. exponential mean one random variables. Now note that by construction, there is a path of length $D_n(1,3) := T_1 + T_{13}$ between vertices 1
and 3 and similarly of length $D_n(2, 3) := T_{12} + T_{23}$ between vertices 2 and 3. Thus, by (4.23) and (4.28)

$$d_n(1, 3) - \log n \leq T_{13} - \frac{1}{2} \log n + T_1 - \frac{1}{2} \log n \xrightarrow{w} \log(1/\hat{E}_1) + \log(1/\hat{E}_3) + \log(\hat{E}_{13}),$$

and

$$d_n(2, 3) - \log n \leq T_{23} - \frac{1}{2} \log n + T_1 - \frac{1}{2} \log n \xrightarrow{w} \log(1/\hat{E}_1) + \log(\hat{E}_{23}/\hat{E}_{12}) + \log(1/\hat{E}_2) + \log(\hat{E}_{12})$$

$$= \log(1/\hat{E}_2) + \log(1/\hat{E}_3) + \log(\hat{E}_{23}),$$

Thus the limiting array $D(3)$ in (4.18) is a limiting upper bound in the weak sense for the array $d_n(3) := (d_n(\alpha, \beta) - \log n : 1 \leq \alpha < \beta \leq 3).$ However, we have equality for $m = 2$ by (4.24). Thus the marginals of $d_n(3)$ converge to the marginals of $D$ as $n \to \infty.$ This implies $d_n(3) \xrightarrow{w} D(3)$ as $n \to \infty.$

This entire construction extends inductively for higher values of $m.$ For example, for $m = 4,$ we consider the random sets $C_1 = \text{SWG}_{12}, C_2 = \text{SWG}_{12}^{(2)}$ and $C_3 = \text{SWG}_{123}^{(3)}$, where $T_{(12)3} = \max(T_{12}, T_{13}).$ Then we define $\mathcal{A}$ as the union of the above three sets and proceed with the construction as above. This completes the proof.

Remark. We learned about this reduction from the sums of collision times to lengths of optimal paths via stochastic domination from [23].

The following is an easy corollary of the proof of the above result. Recall that for any two vertices $\alpha, \beta \in [n],$ $\pi(\alpha, \beta)$ denotes the unique shortest path (geodesic) between them.

Corollary 4.5. Consider the random metric space $\mathcal{S}_n = (\mathcal{K}_n, \{E_i\}_{i \in E_n}).$ Fix $m \geq 2.$

(a) Let $D_n$ be the event that $\exists \alpha \neq \beta \neq \gamma \in [m]$ such that $\gamma \pi (\alpha, \beta).$ Then $\mathbb{P}(D_n) \to 0$ as $n \to \infty.$

(b) Fix $1/2 < \vartheta < 1.$ Consider the smallest-weight graphs $\{\text{SWG}_{\vartheta \log n, i \in [m]}(\alpha, \beta)\}$ from these $m$ vertices at time $\vartheta \log n.$ Then w.h.p. the shortest paths $\pi(\alpha, \beta)$ are contained in the union of these balls, that is, as $n \to \infty,$

$$\mathbb{P}(\pi(\alpha, \beta) \subseteq \bigcup_{i=1}^m \text{SWG}_{\vartheta \log n, i \in [m]}(\alpha, \beta)) \to 1.$$ 

Proof. Part (a) follows from extending (4.21) and (4.25) to general $m.$ Part (b) follows from the above proof, which proves that for any pair of vertices $\alpha, \beta,$ $\pi(\alpha, \beta)$ can be found in $\text{SWG}_{r_n}^{(\vartheta)} \cup \text{SWG}_{r_n}^{( \beta)}$ where $r_n = \frac{1}{2} \log n + O(1).$ 

4.5. Distances between vertices with large minimal edge weight

Fix $\alpha \in \mathbb{R}.$ Recall that

$$N_n(\alpha) = \sum_{i=1}^n 1\{X_i \geq \log n - \alpha\}$$
denotes the number of vertices with minimum outgoing edge length at least \( \log n - \alpha \). Fix \( m \geq 2 \) and condition on the event \( N_n(\alpha) = m \). Let \( V_1, \ldots, V_m \) denote the \( m \) vertices for which \( X_{(V_j)} \geq \log n - \alpha \).

Our aim in this section is to understand, conditionally on the event \( \{N_n(\alpha) = m\} \), the asymptotic joint distribution of \( (d_w(V_i, V_j) : i < j \in [m]) \). Recall the array \( D(m) \) from (4.18) giving the asymptotic joint distribution of the re-centred (by \( \log n \)) length of smallest paths between \( m \) typical vertices in \( \mathcal{S}_n \). The main aim of this section is to prove the following result.

**Proposition 4.6 (distances between vertices with large minimal edge weight).** Fix \( \alpha \in \mathbb{R} \) and \( m \geq 2 \). Recall (4.18). Conditionally on \( N_n(\alpha) = m \), as \( n \to \infty \),

\[
(d_w(V_i, V_j) - 3 \log n + 2\alpha)_{i,j \in [m], i < j} \overset{w}{\to} (\Lambda_i + \Lambda_j - \Lambda_{ij})_{i,j \in [m], i < j} = D(m). \tag{4.31}
\]

We start with two preparatory results, Lemmas 4.7 and 4.8, which we then use to complete the proof of the proposition at the end of this section. Let us start by disentangling exactly what the conditioning event \( \{N_n(\alpha) = m\} \) implies about the edge length distribution. To ease notation, we assume without loss of generality that \( V_1 = i \). Then this conditioning implies that the edge length distributions can be formulated in the following two rules.

(a) **Translation.** Every edge \( E_e' \) incident to one of the vertices in \( [m] \) is conditioned to be at least \( \log n - \alpha \). By the memoryless property of the exponential distribution, we can write \( E_e' = \log n - \alpha + E_e \), where \( (E_e) \) is an independent family of mean \( n \) independent exponential random variables.

(b) **Conditioning.** For every vertex \( i \notin [n] \setminus [m] \), the edges \( (E_{i,j}')_{j \in [m]} \) are independent exponential mean \( n \) random variables conditioned on

\[
X_{(i,[m+1:n])} := \min_{m+1 \leq j \leq n} E_{i,j}' < \log n - \alpha. \tag{4.32}
\]

Let us use our original metric space \( \mathcal{S}_n \) to sequentially overlay the effect of the above two effects. Recall that we have used \( \pi(i,j) \) for the smallest-weight path between \( i,j \) in \( \mathcal{S}_n \). The following lemma deals with the effect of the simpler translation event (without dealing with the conditioning), and will be the starting point of our analysis.

**Lemma 4.7.** Fix \( m \geq 1 \) and consider the metric space \( \mathcal{S}_n \). For every edge \( e \) incident to one of the vertices in \( [m] \), replace the edge \( E_e \) by \( E_e + \log n - \alpha \). Leave all other edges unchanged. Call this new metric space \( \mathcal{S}_n'(tr) \). Write \( \pi'(i,j) \) for the smallest-weight path between \( i,j \) and write \( d_{w}'(i,j) \) for the corresponding metric. Then, for all \( i,j \in [m] \), and on the event that \( E_{ij} + \log n - \alpha > d_{w}(i,j) \) for all \( i,j \in [m] \),

\[
\pi'(i,j) = \pi(i,j), \quad d_{w}'(i,j) = d_{w}(i,j) + 2\log n - 2\alpha. \tag{4.33}
\]

In particular,

\[
(d_{w}'(i,j) - 3 \log n + 2\alpha)_{i,j \in [m], i < j} \overset{w}{\to} (\Lambda_i + \Lambda_j - \Lambda_{ij})_{i,j \in [m], i < j}. \tag{4.34}
\]

**Proof.** The distributional convergence follows from (4.33) and Proposition 4.4. Equation (4.33) follows since we can construct the smallest-weight path problem for \( \mathcal{S}_n'(tr) \) restricted to \( [m] \) as
follows. To $\mathcal{S}_n$ adjoin $m$ new vertices \{i′ : i′ ∈ [m]\}. Each new vertex i′ has only one edge, namely, to vertex i of length $\log n - \alpha$. Call this new metric space $\mathcal{S}_n'$ and the corresponding metric $d_{rw}'$ and smallest-weight path $\pi'(\cdot, \cdot)$. Then the metric space $\mathcal{S}_n'(\text{tr})$ can be constructed as follows. For $i, j \in [m]$ let $d_{rw}'(i, j) = d_{rw}'(i', j')$ and $\pi'(i', j') = \{i' \sim i\} \cup \pi'(i, j) \cup \{j \sim j'\}$.

By the translation rule (a) and conditionally on $N_n(\alpha) = m$, the distances $(d_{uw}(V_i, V_j))_{i, j \in [m], i < j}$ are identical in distribution (assuming w.h.p.) to $(d_{rw}'(i, j) - 3\log n + 2\alpha)_{i, j \in [m], i < j}$ conditionally on $N_{n,m}(\alpha) = 0$, where $N_{n,m}(\alpha)$ denotes the number of vertices in $[n] \setminus [m]$ for which $X_{(i),\lceil m+1, n]} \geq \log n - \alpha$, where $X_{(i),\lceil m+1, n]}$ is defined in (4.32). In order to transfer (4.34) in Lemma 4.7 to the related statement (4.31) in Proposition 4.6, it suffices to prove that the distance matrix $(d_{rw}'(i, j) - \log n)_{i, j \in [m], i < j}$ is asymptotically independent of $N_{n,m}(\alpha)$. That is the content of the next lemma.

**Lemma 4.8.** Fix $\alpha \in \mathbb{R}$ and $m \geq 1$. The matrix of distances $(d_{rw}'(i, j) - \log n)_{i, j \in [m], i < j}$ is asymptotically independent of $N_{n,m}(\alpha)$. Consequently, conditionally on $N_{n,m}(\alpha) = 0$,

$$
(d_{uw}(i, j) - \log n)_{i, j \in [m], i < j} \xrightarrow{\text{w}} (\Lambda_i + \Lambda_j - \Lambda_{ij})_{i, j \in [m], i < j}.
$$

**Proof.** Fix $\ell \geq 0$ and a matrix of weights $(w_{ij})_{i, j \in [m], i < j}$. The statement of the lemma is equivalent to the statement that, for every $(w_{ij})_{i, j \in [m], i < j} \in \mathbb{R}^{m(m-1)/2}$ and $m \geq 1, \ell \geq 0$,

$$
\mathbb{P}\left((d_{uw}(i, j) - \log n)_{i, j \in [m], i < j} \leq (w_{ij})_{i, j \in [m], i < j}, N_{n,m}(\alpha) \leq \ell\right) = \mathbb{P}\left((d_{uw}(i, j) - \log n)_{i, j \in [m], i < j} \leq (w_{ij})_{i, j \in [m], i < j}\right) \mathbb{P}(N_{n,m}(\alpha) \leq \ell) + o(1).
$$

We prove this by proving both inequalities. We note that $(d_{uw}(i, j) - \log n)_{i, j \in [m], i < j}$, as well as $N_{n,m}(\alpha)$, are (sequences of) increasing random variables in the edge weights $(E_{x})_{x}$. Thus, by the FKG-inequality for continuous random variables as stated in [19, Section 6], the left-hand side of (4.36) is always at least the first term on the right-hand side of (4.36). Note that the condition [19, inequality (6.6)] holds with equality since we deal with independent random variables. This establishes the lower bound (even without the error term).

We next prove the upper bound. We split the edges into two disjoint sets. Let $E_1 = \{e : \mathbb{E}_x > \log n - \alpha\}$ and $E_2 = \{e : \mathbb{E}_x \leq \log n - \alpha\}$. Clearly, $N_{n,m}(\alpha)$ is determined by the edges in $E_1$, since $N_{n,m}(\alpha)$ is the number of vertices in $[n] \setminus [m]$ for which every edge incident to it is in $E_1$. Let $d_{uw}(i, j)$ be the distance between $i$ and $j$ only using edges in $E_2$, so that $d_{uw}(i, j) \geq d_{uw}(i, j)$. We bound from above

$$
\mathbb{P}\left((d_{uw}(i, j) - \log n)_{i, j \in [m], i < j} \leq (w_{ij})_{i, j \in [m], i < j}, N_{n,m}(\alpha) \leq \ell\right)
$$

$$
\leq \mathbb{P}\left((d_{uw}(i, j) - \log n)_{i, j \in [m], i < j} \leq (w_{ij})_{i, j \in [m], i < j}, N_{n,m}(\alpha) \leq \ell\right) + \mathbb{P}\left((d_{uw}(i, j))_{i, j \in [m], i < j} \neq (d_{uw}(i, j))_{i, j \in [m], i < j}\right),
$$

and bound both terms from above. In the first term, we note that $(d_{uw}(i, j) - \log n)_{i, j \in [m], i < j} \leq (w_{ij})_{i, j \in [m], i < j}$
Lemma 4.7 with the asymptotic independence of

\begin{equation}
\mathbb{P}\left((\bar{d}_w(i,j) - \log n)_{i,j \in [m], i < j} \leq (w_{ij})_{i,j \in [m], i < j}, N_{n,m}(\alpha) \leq \ell\right) \leq \mathbb{P}\left(\{(\bar{d}_w(i,j) - \log n)_{i,j \in [m], i < j} \leq (w_{ij})_{i,j \in [m], i < j}\} \cap \{N_{n,m}(\alpha) \leq \ell\}\right).
\end{equation}

By the BK-inequality applied to continuous random variables (see [16, Theorem 1.2]),

\begin{equation}
\mathbb{P}\left((\bar{d}_w(i,j) - \log n)_{i,j \in [m], i < j} \leq (w_{ij})_{i,j \in [m], i < j}, N_{n,m}(\alpha) \leq \ell\right) \leq \mathbb{P}\left((\bar{d}_w(i,j) - \log n)_{i,j \in [m], i < j} \leq (w_{ij})_{i,j \in [m], i < j}\right) \mathbb{P}(N_{n,m}(\alpha) \leq \ell).
\end{equation}

By the fact that \(\bar{d}_w(i,j) \geq d_w(i,j)\), the right-hand side of (4.39) is at most the product of probabilities on the right-hand side of (4.36). It remains to show that the second term in (4.37) is \(o(1)\).

By symmetry and the union bound,

\begin{equation}
\mathbb{P}\left((d_w(i,j))_{i,j \in [m], i < j} \neq (\bar{d}_w(i,j))_{i,j \in [m], i < j}\right) \leq m^2 \mathbb{P}(d_w(1,2) \neq \bar{d}_w(1,2)).
\end{equation}

By Proposition 4.4, as well as its proof (see Corollary 4.5(b)), \(\mathbb{P}(d_w(1,2) \neq \bar{d}_w(1,2)) = o(1)\), which completes the proof of the asymptotic independence. To prove (4.35), it suffices to show that \(\liminf_{n \to \infty} \mathbb{P}(N_{n,m}(\alpha) = 0) > 0\), which follows from Proposition 4.3 together with the fact that \(N_{n,m}(\alpha) = N_n(\alpha)\) w.h.p.

**Proof of Proposition 4.6.** By the translation rule (a) and the conditioning rule (b), the distribution of \((d_w(V_i,V_j))_{i,j \in [m], i < j}\) conditionally on \(N_n(\alpha) = m\) is the same as that of \((d_w(i,j))_{i,j \in [m], i < j}\) in \(\mathcal{G}_w(\text{tr})\) conditionally on \(N_{n,m}(\alpha) = 0\). The event that \(E_{ij} + \log n - \alpha > d_w(i,j)\) for all \(i,j \in [m]\) occurs w.h.p., so the identity in (4.33) holds w.h.p. Combining (4.34) in Lemma 4.7 with the asymptotic independence of \((d_w(i,j))_{i,j \in [m], i < j}\) and \(N_{n,m}(\alpha)\) in Lemma 4.8 and the fact that \(N_{n,m}(\alpha) \xrightarrow{w} N(\alpha)\) by Proposition 4.3 completes the proof of Proposition 4.6.

### 4.6. Reduction to distances between vertices with large minimal edge weights

The previous section analysed distances between the vertices whose minimal outgoing edge is large (e.g. \(\log n + O_n(1)\)). The distances between these vertices are then close to \(3 \log n + O_n(1)\). The aim of this section is to show that these are the only vertices that matter for the weight diameter. We achieve this by considering distances between vertices whose minimal outgoing edge is ‘small’ and showing that the distance between such vertices are not large enough to create the diameter and thus can be ignored.

We start with some notation. Fix \(\alpha > 0\) and define

\begin{equation}
R_n(\alpha) = \#\{i,j \in [n]: X_{(i)} \leq \log n - \alpha, X_{(j)} \leq \log n + \alpha/2, d_w(i,j) \geq 3 \log n - \alpha/8\}.
\end{equation}

The random variable \(R_n(\alpha)\) counts the number of ordered pairs of vertices \((i,j) \in [n] \times [n]\) that satisfy that the minimal outgoing edge of vertex \(i\) is less than \(\log n - \alpha\), the minimal outgoing
edge of \( j \) is less than \( \log n + \alpha/2 \) and yet the distance between \( i, j \) is greater than \( 3\log n - \alpha/8 \). The following result gives an upper bound on the expected value of \( R_n(\alpha) \).

**Proposition 4.9 (distances from vertices with small minimal weight).** There exists a constant \( C > 0 \) such that, for all \( \alpha > 0 \),

\[
\limsup_{n \to \infty} \mathbb{E}[R_n(\alpha)] \leq Ce^{-\alpha/16}. \tag{4.42}
\]

**Proof.** By the union bound,

\[
\mathbb{E}[R_n(\alpha)] \leq n^2 \mathbb{P}(d_w(1, 2) \geq 3\log n - \alpha/8, X_{(1)} \leq \log n - \alpha, X_{(2)} \leq \log n + \alpha/2). \tag{4.43}
\]

Note that

\[
(X_{(1)}, X_{(2)}) = \left( \min \left[ \frac{n}{n-2} E_1^*, nE_{12}^* \right], \min \left[ \frac{n}{n-2} E_2^*, nE_{12}^* \right] \right),
\]

where \( E_1^*, E_2^*, E_{12}^* \) are independent exponential random variables with mean 1. Here \( nE_{12}^* \) represents the weight of the direct edge between vertices 1, 2, while for \( i \in \{1, 2\} \), \( nE_i^*/(n-2) \) represents the minimal outgoing edges from vertex \( i \) to the remaining vertices \( [n] \setminus \{1, 2\} \).

On the event \( \{d_w(1, 2) \geq 3\log n - \alpha/8\} \), we have that \( nE_{12}^* \geq d_w(1, 2) \geq 3\log n - \alpha/8 \). As a result, when \( d_w(1, 2) \geq 3\log n - \alpha/8 \), unless

\[
\max \left( \frac{n}{n-2} E_1^*, \frac{n}{n-2} E_2^* \right) > 3\log n - \alpha/8,
\]

we have that

\[
(X_{(1)}, X_{(2)}) = \left( \frac{n}{n-2} E_1^*, \frac{n}{n-2} E_2^* \right). \tag{4.45}
\]

The probability of the event in (4.44) is bounded by \( 3e^{\alpha/8}/n^3 \) for large \( n \). Since \( n^2 e^{\alpha/8}/n^3 \to 0 \), we can ignore the contribution of this in the proof of Proposition 4.9 and assume (4.45).

Let \( V_1 \) be the closest vertex to 1, at distance \( X_{(1)} \) (respectively \( V_2 \) at distance \( X_{(2)} \) from vertex 2). The rest of the smallest-weight path has the same distribution as the smallest-weight path between two sets \( A = \{1, V_1\} \) and \( B = \{2, V_2\} \) in \( \mathcal{S}_n \). Lemma 4.1 thus implies that

\[
d_w(i, j) = X_{(1)} + X_{(2)} + \sum_{k=2}^{N-1} \frac{nE_k^*}{k(n-k)}, \tag{4.46}
\]

where \( N = N_1 \land N_2 \) and \( (N_1, N_2) \) is a uniform pair of distinct vertices from \( [n] \setminus \{1, 2\} \) and \( (E_k^*)_k \geq 1 \) are mean one exponential random variables. The distribution of \( N \) in Lemma 4.1 was given as the number of draws required to draw the first black ball in an urn containing \( N - 4 \) white and 2 black balls; it is easy to check that this is equivalent to the description above. Writing

\[
S_N = \sum_{k=2}^{N-1} \frac{nE_k^*}{k(n-k)},
\]

we get

\[
\mathbb{E}[R_n(\alpha)] \leq n^2 \mathbb{P}(S_N \geq 3\log n - X_{(1)} - X_{(2)} - \alpha/8, X_{(1)} \leq \log n - \alpha, X_{(2)} \leq \log n + \alpha/2). \tag{4.47}
\]
Thus,

\[ E[R_n(\alpha)] \leq n^2 \int_0^{\log n - \alpha} \int_0^{\log n + \alpha/2} e^{-(x+y)(n-2)/n} P(S_N \geq 3 \log n - x - y - \alpha/8) \, dx \, dy. \quad (4.48) \]

To complete the proof, we study the tail behaviour of the random variable \( S_N \).

**Lemma 4.10 (tail behaviour for random sums).** For any constant \( a < 2 \), there exists a constant \( C = C_a < \infty \) such that, for every \( x \geq 0 \),

\[ P(S_N \geq \log n + x) \leq Ce^{-ax}. \quad (4.49) \]

Assuming the lemma for the time being, we use this to complete the proof of Proposition 4.9. Using Lemma 4.10 with \( a = 3/2 \) gives

\[ E[R_n(\alpha)] \leq Cn^2 \int_0^{\log n - \alpha} \int_0^{\log n + \alpha/2} e^{-(x+y)} e^{-a(2 \log n - x - y - \alpha/8)} \, dx \, dy \]

\[ = Cn^{2(1-a)} \int_0^{\log n - \alpha} \int_0^{\log n + \alpha/2} e^{(a-1)(x+y)} e^{a \alpha/8} \, dx \, dy \]

\[ = Ce^{-(a-1)\alpha/2 + a\alpha/8} \leq Ce^{-\alpha/16}. \]

This completes the proof of Proposition 4.9. \( \square \)

**Proof of Lemma 4.10.** We compute the moment generating function of \( S_N \) as

\[ M_{S_N}(a) = \sum_{j=2}^{n-2} \mathbb{P}(N = j) \mathbb{E}[e^{aS_j}] = \sum_{j=2}^{n-2} \mathbb{P}(N = j) \prod_{k=2}^{j-1} \frac{k(n-k)}{k(n-k) - an} \]

\[ = \sum_{j=2}^{n-2} \mathbb{P}(N = j) \exp \left( - \sum_{k=2}^{j-1} \log \left( 1 - \frac{an}{k(n-k)} \right) \right). \quad (4.51) \]

Thus,

\[ P(S_N \geq \log n + x) \leq e^{-a(\log n + x)} M_{S_N}(a) \]

\[ \leq e^{-a(\log n + x)} \sum_{j=2}^{n-2} \mathbb{P}(N = j) \exp \left( - \sum_{k=2}^{j-1} \log \left( 1 - \frac{an}{k(n-k)} \right) \right). \quad (4.52) \]

Using \( a < 2 \), we have \( an/[k(n-k)] < 1 \), since \( k, n-k \geq 2 \). Therefore, via Taylor expansion,

\[ \log \left( 1 - \frac{an}{k(n-k)} \right) \leq \frac{an}{k(n-k)} + O\left( \frac{n^2}{[k(n-k)]^2} \right), \quad (4.53) \]

Using

\[ \frac{n}{k(n-k)} = \frac{1}{k} + \frac{1}{n-k}, \]
we arrive at
\[ P(S_N \geq \log n + x) \leq e^{-a(\log n + x)}M_{S_N}(a) \]
\[ \leq Ce^{-a(\log n + x)} \sum_{j=2}^{n-2} P(N = j) \exp \left( a \sum_{k=2}^{j-1} \left[ \frac{1}{k} + \frac{1}{n-k} \right] \right) \]
\[ \leq Ce^{-ax} \sum_{j=2}^{n-2} P(N = j) e^{a[\log(j/n) - \log(1-j/n)]} \]
\[ = Ce^{-ax} \mathbb{E} \left( \left( \frac{N/n}{1-N/n} \right)^a \right). \]

Note that
\[ P(N = j) = 2(n-j)/(n-2)(n-3), \]
so that, by dominated convergence,
\[ \mathbb{E} \left( \left( \frac{N/n}{1-N/n} \right)^a \right) = \sum_{j=2}^{n-2} 2(n-j)/(n-2)(n-3) \left( \frac{j/n}{1-j/n} \right)^a \to \int_0^1 \frac{u^a}{(1-u)^2} 2(1-u) \, du < \infty, \quad (4.54) \]
whenever \( a < 2. \)

4.7. The limiting random variable
In this section we prove the finiteness of the random variable \( \Xi = \max_{s<t}(Y_s + Y_t - \Lambda_{st}) \) in (2.3), which Theorem 2.1 asserts is the limit of the re-centred diameter. In the following lemma we give an alternative expression for its distribution.

Lemma 4.11 (the limiting random variable). Let \( Q = e^{-\Xi}. \) Then
\[ Q = \min_{s<t} \frac{S_i S_j}{E'_{st}}, \quad (4.55) \]
where \( S_i = \sum_{l=1}^i E'_i \) and \( (E'_i)_{i \geq 1} \) and \( (E'_{st})_{s<t} \) are i.i.d. exponential random variables with mean 1. In particular, for every \( x > 0, \)
\[ P(Q > x) = \mathbb{E} \left[ \prod_{1 \leq s < t} l(1 - e^{-S_i / x}) \right], \quad (4.56) \]
and \( P(Q > x) \in (0, 1) \) for every \( x > 0. \) Further, \( \lim_{x \downarrow 0} P(Q > x) = 1. \)

Proof. We note that we can write \( -\Lambda_{st} = \log(E'_{st}) \) and \( Y_s = -\log(S_i). \) Indeed, the point process \( (e^{-Y_s})_{s \geq 1} \) is a standard Poisson process. Thus,
\[ e^{-\Xi} = \min_{s<t} e^{\log(S_i) + \log(S_j) - \log(E'_{st})} = Q. \quad (4.57) \]
Equation (4.56) immediately follows. To prove that \( P(Q > x) \in (0, 1) \) for every \( x > 0, \) we note that \( P(Q > x) < 1 \) follows immediately from (4.56) since each of the terms in the product is \( < 1 \)
a.s. To show that \( P(Q > x) > 0 \), we first note that
\[
P(Q > x) \geq E \left[ \prod_{1 \leq s < t} (1 - e^{-S_s / x}) \mathbb{1}_{\{S_1 > 1\}} \right]
\]
\[
= E \left[ \prod_{1 \leq s < t} (1 - e^{-S_s / x}) \mid S_1 > 1 \right] P(S_1 > 1). \tag{4.58}
\]

We compute that \( P(S_1 > 1) = 1/e \), and observe that by the memoryless property of the exponential random variable \( S_1 \), conditionally on \( S_1 > 1 \), the distribution of \( (S_t)_{t \geq 1} \) is equal to \( (S_t + 1)_{t \geq 1} \). Thus,
\[
P(Q > x) \geq e^{-1} E \left[ \prod_{1 \leq s < t} (1 - e^{-(S_s + 1) / (S_t + 1)}) \right]
\]
\[
\geq e^{-1} \exp \left( \sum_{1 \leq s < t} E \left[ \log(1 - e^{-(S_s + 1) / (S_t + 1)}) \right] \right). \tag{4.59}
\]

Next, we compute, using Fubini,
\[
\sum_{1 \leq s < t} E \left[ \log(1 - e^{-(S_s + 1) / (S_t + 1)}) \right]
\]
\[
= \sum_{1 \leq s < t} \int_0^\infty du \int_0^\infty dv \frac{u^{s-1}}{(s-1)!} \frac{v^{t-s-1}}{(t-s-1)!} e^{-(u+v) \log(1 - e^{-(u+1) / (u+v+1)})}
\]
\[
= \int_0^\infty du \int_0^\infty dv \sum_{1 \leq s < t} \frac{u^{s-1}}{(s-1)!} \frac{v^{t-s-1}}{(t-s-1)!} e^{-(u+v) \log(1 - e^{-(u+1) / (u+v+1)})}
\]
\[
= \int_0^\infty \int_0^\infty \log(1 - e^{-(u+1) / (u+v+1)}) \, du \, dv < \infty.
\]

This proves that \( P(Q > x) > 0 \) for every \( x > 0 \). Similarly, for any \( \varepsilon > 0 \), by conditioning on \( S_1 \geq \varepsilon \),
\[
P(Q > x) \geq e^{-\varepsilon} E \left[ \prod_{1 \leq s < t} (1 - e^{-(S_s + \varepsilon) / (S_t + \varepsilon)}) \right]
\]
\[
\geq e^{-\varepsilon} \exp \left( \sum_{1 \leq s < t} E \left[ \log(1 - e^{-(S_s + \varepsilon) / (S_t + \varepsilon)}) \right] \right)
\]
\[
= e^{-\varepsilon} \exp \left( \int_0^\infty \int_0^\infty \log(1 - e^{-(u+1) / (u+v+1)}) \, du \, dv \right)
\]
\[
= e^{-\varepsilon} \exp \left( e^{-2} \int_0^\infty \int_0^\infty \log(1 - e^{-(u+1) / (u+v+1)}) \, du \, dv \right),
\]
where the last equality follows by substituting \((u', v') = (\varepsilon u, \varepsilon v)\). For every \( \varepsilon > 0 \), as \( x \searrow 0 \), the integrand converges pointwise to zero and thus, by monotone convergence, the integral also converges to 0. Therefore, for every \( \varepsilon > 0 \),
\[
\liminf_{x \searrow 0} P(Q > x) \geq e^{-\varepsilon}, \tag{4.62}
\]
and letting \( \varepsilon \searrow 0 \), this completes the proof that \( \lim_{x \searrow 0} P(Q > x) = 1 \).\[\square\]
4.8. The limiting maximization problem
In this section we combine the various ingredients proved in the previous sections to prove the
distributional convergence in Theorem 2.1. We defer the proof of the convergence of moments to
the next section. By Proposition 4.3, first note that given any \( \varepsilon > 0 \), one can choose \( \alpha \) large but
finite (independent of \( n \)) such that for all \( n \) large, \( \Pr(N_n(\alpha) \geq 2) \geq 1 - \varepsilon \). Thus by Proposition 4.6,
with probability greater than \( 1 - \varepsilon \) for all large \( n \),
\[
\text{Diam}_w(K_n) - 3 \log n \geq d_w(V_1, V_2) - 3 \log n \xrightarrow{w} -2\alpha + \Lambda_1 + \Lambda_2 - \Lambda_{12}. \tag{4.63}
\]
As a result the sequence \( \{ (\text{Diam}_w(K_n) - 3 \log n) : n \geq 1 \} \) is tight. Further, using Proposition 4.9,
given \( \varepsilon > 0 \), one can choose \( \alpha = \alpha(\varepsilon) \) large enough such that with probability \( \geq 1 - \varepsilon \) for all
large \( n \),
\[
\text{Diam}_w(K_n) = \max_{s < t \leq N_n(\alpha)} d_w(V_s, V_t). \tag{4.64}
\]
We note that, again using Propositions 4.6 and 4.3,
\[
\max_{s < t \leq N_n(\alpha)} d_w(V_s, V_t) - 3 \log n \xrightarrow{w} \max_{s < t \leq N(\alpha)} (\Lambda_s + \Lambda_t - \Lambda_{st} - 2\alpha), \tag{4.65}
\]
where \( N(\alpha) \) is a Poisson random variable with mean \( e^\alpha \) and the Gumbel variables are independ-
et of \( N(\alpha) \). As a result,
\[
\text{Diam}_w(K_n) - 3 \log n \xrightarrow{w} \Xi^*, \tag{4.66}
\]
where \( \Xi^* \) is the distributional limit as \( \alpha \to \infty \) of the right-hand side of (4.65), that is,
\[
\max_{s < t \leq N(\alpha)} (\Lambda_s + \Lambda_t - \Lambda_{st} - 2\alpha) \xrightarrow{w} \Xi^*. \tag{4.67}
\]
Theorem 4.12 below shows that this weak limit exists and that \( \Xi^* = \Xi \) defined in (2.3). This
completes the proof of the first assertion of the main theorem, thus showing that
\[
\text{Diam}_w(K_n) - 3 \log n \xrightarrow{w} \Xi. \tag{4.68}
\]

**Theorem 4.12 (the limiting variable \( \Xi \)).** As \( \alpha \to \infty \),
\[
\max_{s < t \leq N(\alpha)} (\Lambda_s + \Lambda_t - \Lambda_{st} - 2\alpha) \xrightarrow{w} \Xi, \tag{4.68}
\]
where \( \Xi \) is defined in (2.3).

**Proof.** First note that as \( \alpha \to \infty \),
\[
e^{-\alpha N(\alpha)} \xrightarrow{P} 1. \tag{4.69}
\]
For each fixed \( \alpha > 0 \), define the random variables
\[
\hat{\Xi}_\alpha := \max_{s < t \leq N(\alpha)} (\Lambda_s + \Lambda_t - \Lambda_{st} - 2\alpha), \quad \Xi_\alpha := \max_{s < t \leq e^\alpha} (\Lambda_s + \Lambda_t - \Lambda_{st} - 2\alpha).
\]
It then suffices to prove that
\[
\Xi_\alpha := \max_{s < t \leq e^\alpha} (\Lambda_s + \Lambda_t - \Lambda_{st} - 2\alpha) \xrightarrow{w} \Xi. \tag{4.70}
\]
Then, for each \( \varepsilon > 0 \), one can sandwich \( \hat{\Xi}_\alpha \) between \( \Xi_{\alpha - \varepsilon} \) and \( \Xi_{\alpha + \varepsilon} \) (with high probability as \( \alpha \to \infty \) by (4.69)), and then first let \( \alpha \to \infty \) and then \( \varepsilon \to 0 \) to get the corresponding result for \( \hat{\Xi}_\alpha \).

Recall from Section 2, that the Poisson point process \( \mathcal{P} = (Y_s)_{s \geq 1} \) with intensity measure given by the density function \( \lambda(y) = e^{-y} \). Also recall from (2.3) that we have defined \( \Xi \) as

\[
\Xi := \max_{s < t} (Y_s + Y_t - \Lambda_{st}).
\]

For any fixed \( A > 0 \), let \( \mathcal{P}(A) \) denote \( \mathcal{P} \) restricted to the interval \([-A, \infty)\). Write

\[
\Xi(A) := \max_{s < t: Y_s, Y_t \in \mathcal{P}(A)} (Y_s + Y_t - \Lambda_{st}).
\]

Thus, \( \Xi(A) \) is the maximum of corresponding pairs \((s, t)\) whose point process values satisfy \( Y_s, Y_t > -A \). Intuitively, one would expect that \( \Xi = \Xi(A) \) for large \( A \). We now make his intuition precise. Define

\[
\mathcal{R}^{(1)}(A) := \max_{s < t: Y_s, Y_t \leq -A} (Y_s + Y_t - \Lambda_{st}),
\]

and, for \( A < B \), let

\[
\mathcal{R}^{(2)}(A, B) := \max_{s < t: Y_s \geq A, Y_t \leq -(A+B)} (Y_s + Y_t - \Lambda_{st}).
\]

The random variable \( \mathcal{R}^{(1)}(A) \) is the supremum between pairs \((s, t)\) such that \( Y_s, Y_t \leq -A \) while \( \mathcal{R}^{(2)}(A, B) \) corresponds to the supremum between pairs of points \((s, t)\) such that \( Y_s > -A \) but \( Y_t < -(A + B) \). Note that, for any \( z \),

\[
\{ \Xi = \Xi(A + B) \} \supseteq \{ \Xi(A) > z, \mathcal{R}^{(1)}(A) < z, \mathcal{R}^{(2)}(A, B) < z \}. \tag{4.71}
\]

Consider the point process

\[
\mathcal{P}_\alpha^* = \sum_{\alpha} \mathcal{P} \delta\{\Lambda_{s} - \alpha\}.
\]

When arranged in increasing order, write this point process as \( Y_1(\alpha) > Y_2(\alpha) > \cdots \). Standard extreme value theory implies that

\[
\mathcal{P}_\alpha^* \xrightarrow{w} \mathcal{P} \quad \text{as } \alpha \to \infty, \tag{4.72}
\]

where \( \xrightarrow{w} \) denotes convergence in distribution in the space of point measures on \( \mathbb{R} \) equipped with the vague topology. Analogously to \( \Xi(A), \mathcal{R}^{(1)}(A), \mathcal{R}^{(2)}(A, B) \), we define the random variables \( \Xi_\alpha(A), \mathcal{R}^{(1)}_\alpha(A), \mathcal{R}^{(2)}_\alpha(A, B) \), that is,

\[
\Xi_\alpha(A) := \max_{s < t: Y_s(\alpha), Y_t(\alpha) \in \mathcal{P}_\alpha(A)} (Y_s(\alpha) + Y_t(\alpha) - \Lambda_{st}),
\]

where \( \mathcal{P}_\alpha(A) \) is the point process \( \mathcal{P}_\alpha \) restricted to the interval \([-A, \infty)\). We define \( \mathcal{R}^{(1)}_\alpha(A), \mathcal{R}^{(2)}_\alpha(A) \).
similarly. As before, for any \( z \),
\[
\{ \Xi_\alpha(A+B) \} \supseteq \{ \Xi_\alpha(A) > z, R^{(1)}_\alpha(A) < z, R^{(2)}_\alpha(A,B) < z \}.
\] (4.73)

The weak convergence in (4.72) immediately implies that, for any fixed \( A \),
\[
\Xi_\alpha(A) \xrightarrow{w} \Xi(A) \quad \text{as} \quad \alpha \to \infty.
\] (4.74)

The following lemma formalizes the notion that for large \( A \), \( \Xi = \Xi(A) \) and, similarly, when \( \alpha \) is large, \( \Xi_\alpha(A) = \Xi_\alpha \). Let us give an intuitive description of how this is achieved. The argument proceeds by showing that for large \( A \), each of the random variables \( R^{(1)}_\alpha(A) \), \( R^{(2)}_\alpha(A,B) \) take large negative values.

\textbf{Proposition 4.13.}

(a) Fix \( x \in \mathbb{R} \). Then
\[
\limsup_{A \to \infty} \mathbb{P}(R^{(1)}(A) > x) = 0.
\]
Further, for each fixed \( A \),
\[
\limsup_{B \to \infty} \mathbb{P}(R^{(2)}(A,B) > x) = 0.
\]

(b) Fix \( x \in \mathbb{R} \). Then,
\[
\limsup_{A \to \infty} \limsup_{\alpha \to \infty} \mathbb{P}(R^{(1)}_\alpha(A) > x) = 0.
\]
Further, for each fixed \( A \),
\[
\limsup_{B \to \infty} \limsup_{\alpha \to \infty} \mathbb{P}(R^{(2)}_\alpha(A,B) > x) = 0.
\]

We prove the statements in Proposition 4.13 one by one.

\textbf{Proof of Proposition 4.13(a).} We start with \( R^{(1)}(A) \). To simplify notation, we also restrict ourselves to the case \( x = 0 \). The general \( x \) case is identical.

Write
\[
\mathcal{N}^{(1)}(A) := \#\{(s,t) : Y_s, Y_t < -A, Y_s + Y_t - \Lambda_{st} \geq 0 \}.
\]
It is enough to show \( \limsup_{A \to \infty} \mathbb{E}(\mathcal{N}^{(1)}(A)) = 0 \). Conditioning on the point process \( \mathcal{P} \), we get
\[
\mathbb{E}(\mathcal{N}^{(1)}(A) | \mathcal{P}) = \sum_{(s,t), s < t, Y_s < -A} e^{-\beta(Y_s + Y_t)}.
\]
Fix \( a > 1 \). We use the fact that we can choose \( A \) so large such that \( e^{-e^{C+D}} < e^{-aC}e^{-aD} \) for all \( C, D > A \). This leads to
\[
\mathbb{E}(\mathcal{N}^{(1)}(A) | \mathcal{P}) \leq \sum_{(s,t), s < t, Y_s < -A} e^{aY_s}e^{aY_t}.
\]
Since \( \{Y_s \in \mathcal{P} : Y_s \leq -A\} \) is just a Poisson point process on the interval \((-\infty, -A]\) with density \( \lambda(x) = e^{-x} \), properties of Poisson processes [20, Eqn 3.14] implies that, as \( A \to \infty \),

\[
E \left( \sum_{(s,t), s < t, \ Y_s, Y_t < -A} e^{Y_s} e^{Y_t} \right) = \frac{1}{2} \left( \int_{-\infty}^{-A} e^{\alpha x} e^{-x} \, dx \right)^2 = \frac{1}{2} e^{-2(a-1)A} \to 0.
\]

This shows that \( \limsup_{A \to \infty} E(N^{(1)}(A)) = 0 \) and thus completes the proof.

Next fix \( A \) and let us deal with \( \mathcal{R}^{(2)}(A, B) \). Here we use the fact that \( \mathcal{P}(A) \) and \( \mathcal{P}^{c}(A+B) := \mathcal{P} \setminus \mathcal{P}^{c}(A+B) \) are independent Poisson point processes on the sets \([-A, \infty) \) and \((-\infty, -(A+B)] \) with intensity measure with density \( \lambda(y) = e^{-y} \). We condition on \( \mathcal{P}(A) \). Fix a point \( Y_s \) in \( \mathcal{P}(A) \). Then

\[
P \left( \sup_{Y_t < -(A+B)} (Y_s + Y_t - \Lambda_{st}) < z\big| \mathcal{P}(A) \right) = E \left( \prod_{t: Y_t < -(A+B)} \left( 1 - \exp\left( -e^{-(Y_t - (z-Y_s))} \right) \right) \right). \tag{4.75}
\]

Thus to complete the proof, it is enough to show, for any \( z^* \) and \( A \),

\[
\lim_{B \to \infty} E \left( \prod_{t: Y_t < -(A+B)} \left( 1 - \exp\left( -e^{-(Y_t - (z-Y_s))} \right) \right) \right) \to 1. \tag{4.76}
\]

By the dominated convergence theorem, it is enough to show that, as \( B \to \infty \),

\[
\prod_{t: Y_t < -(A+B)} \left( 1 - \exp\left( -e^{-(Y_t - (z-Y_s))} \right) \right) \xrightarrow{P} 1.
\]

Taking logarithms, this is equivalent to showing that, as \( B \to \infty \),

\[
\sum_{t: Y_t < -(A+B)} \log \left( 1 - \exp\left( -e^{-(Y_t - (z-Y_s))} \right) \right) \xrightarrow{P} 0.
\]

In turn, this is equivalent to showing that, as \( B \to \infty \),

\[
\sum_{t: Y_t < -(A+B)} \exp\left( -e^{-(Y_t - (z-Y_s))} \right) \xrightarrow{P} 0.
\]

By Campbell’s theorem [20],

\[
E \left( \sum_{t: Y_t < -(A+B)} \exp\left( -e^{-(Y_t - z^*)} \right) \right) = \int_{-\infty}^{-(A+B)} \exp\left( -e^{-(y - z^*)} \right) e^{-y} \, dy
\]

\[
= e^{z^*} \exp\left( -e^{A+B+z^*} \right) \to 0,
\]

as \( B \to \infty \). This completes the proof of (4.76) and thus of part (a) of Proposition 4.13.

Proof of Proposition 4.13(b). This closely follows the proof of part (a). Without giving a full proof, we mainly highlight the differences. We again start with \( \mathcal{R}^{(3)}_{\alpha}(A) \) and again restrict ourselves to the case \( x = 0 \). The general \( x \) case is identical.

Write

\[
\mathcal{N}^{(3)}_{\alpha}(A) := \# \{ (s, t) : Y_s(\alpha), y_t(\alpha) < -A, Y_s(\alpha) + Y_t(\alpha) - \Lambda_{st} \geq 0 \}.
\]
Recall that \(C_{ij}\).

**4.9. Convergence of moments**

Completing the proof of Theorem 4.12.

It is enough to show that

\[
\limsup_{A \to \infty} \limsup_{\alpha \to \infty} E(\mathcal{N}_\alpha^{(1)}(A)) = 0.
\]

Conditioning on the point process \(\mathcal{P}_\alpha^*\), we now get

\[
E(\mathcal{N}_\alpha^{(1)}(A) | \mathcal{P}_\alpha^*) = \sum_{(s,t), s < t, Y_s(\alpha), Y_t(\alpha) < -A} \exp(-e^{-(Y_s(\alpha) + Y_t(\alpha))})
\]

\[
= \sum_{1 \leq s < t \leq e^\alpha} \mathbb{1}_{\{\Lambda_s, \Lambda_t < -A + \alpha\}} \exp(-e^{-(\Lambda_s - \alpha) - (\Lambda_t - \alpha)}).
\]

Taking expectations and using that \(\Lambda_s, \Lambda_t\) are independent for \(s < t\) leads to

\[
E(\mathcal{N}_\alpha^{(1)}(A)) \leq \int_{-\infty}^{e^{-A + \alpha}} \int_{-\infty}^{e^{-A + \alpha}} e^{-(u - \alpha)} \exp\left(\exp(-e^{-(u - \alpha)})\right) du dv.
\]

This integral can be bounded by

\[
E(\mathcal{N}_\alpha^{(1)}(A)) \leq \int_{-\infty}^{-A} \int_{-\infty}^{-A} e^{-u} e^{-v} \exp(-e^{-u - v}) du dv,
\]

which is independent of \(\alpha\) and converges to 0 as \(A \to \infty\). The proof for \(\mathcal{R}_\alpha^{(2)}(A, B)\) is similar and will be omitted.

**Completing the proof of Theorem 4.12.** Proposition 4.13 shows that given any \(\varepsilon > 0\) one can choose finite \(\alpha = \alpha(\varepsilon)\) and \(A = A(\varepsilon)\) such that \(\mathcal{P}(\Xi = \Xi(A))\) and \(\mathcal{P}(\Xi_\alpha(A) = \Xi_\alpha)\) are both at least \(1 - \varepsilon\). Using (4.74), (4.71) and (4.73) completes the proof of Theorem 4.12.

**4.9. Convergence of moments**

Recall that \(C_{ij} = d_n(i, j)\). We need to show that

\[
E\left[\max_{i,j \in [n]} C_{ij}\right] - 3\log n \to E[\Xi], \quad \text{Var}\left(\max_{i,j \in [n]} C_{ij}\right) \to \text{Var}(\Xi).
\]

Since we have already shown convergence in distribution, by uniform integrability for any \(p \geq 1\), to prove that

\[
E\left[\left(\max_{i,j \in [n]} C_{ij} - 3\log n\right)^p\right] \to E[\Xi^p],
\]

it suffices to prove that, for some integer \(q\) with \(q > p/2\),

\[
E\left[\left(\max_{i,j \in [n]} C_{ij} - 3\log n\right)^{2q}\right] = O(1).
\]

Combined with convergence in distribution, this implies convergence of the moments as well as existence of the moments of the limit random variable \(\Xi\). Note that

\[
E\left[\left(\max_{i,j \in [n]} C_{ij} - 3\log n\right)^{2q}\right] = E\left[\left(\max_{i,j \in [n]} C_{ij} - 3\log n\right)^{2q}\right] + E\left[\left(\max_{i,j \in [n]} C_{ij} - 3\log n\right)^{2q}\right].
\]

We start by analysing the first term on the right-hand side of (4.82) by deriving an upper bound on \(\max_{i,j \in [n]} C_{ij} - 3\log n\), and then prove a lower bound on \(\max_{i,j \in [n]} C_{ij} - 3\log n\) to obtain a bound on the second term on the right-hand side of (4.82).
Upper bound. Let us analyse the first term in (4.82) and show that
\[ \mathbb{E} \left[ \left( \max_{i,j \in [n]} C_{ij} - 3 \log n \right)^{2q} \right] = O(1). \]

To prove this assertion, it is enough to show that there exist \( N, \alpha \) such that, for all large \( n > N \) and \( x \geq \alpha \), the random variable \( \max_{i,j \in [n]} C_{ij} - 3 \log n \) has exponential upper tails in the sense that there exist constants \( \kappa_1, \kappa_2 > 0 \) (independent of \( x \)) such that
\[ \mathbb{P} \left( \max_{i,j \in [n]} C_{ij} - 3 \log n > x \right) \leq \kappa_1 e^{-\kappa_2 x}. \] (4.83)

Now note that
\[ 1 \left\{ \max_{i,j \in [n]} C_{ij} - 3 \log n > x \right\} \leq 1 \left\{ \max_{i \in [n]} X_i > \log n + 4x \right\} + R^{(1)}_n(x) + R^{(2)}_n(x). \] (4.84)

Here \( R^{(1)}_n(x) = R_n(8x) \) as in (4.41), that is,
\[ R^{(1)}_n(x) = \# \{ i, j \in [n] : X_{(i)} \leq \log n - 8x, X_{(j)} \leq \log n + 4x, d_w(i, j) \geq 3 \log n - x \}, \]
while
\[ R^{(2)}_n(x) := \# \{ (i, j) : X_{(i)} > \log n - 8x, X_{(j)} > \log n - 8x, d_w(i, j) > 3 \log n + x \}. \]

Recall that, for any \( \alpha \in \mathbb{R} \), \( N_n(\alpha) \) denotes the number of vertices \( i \) with \( X_{(i)} \geq \log n - \alpha \). For the first term in (4.84), since
\[ \mathbb{P} \left( \max_{i \in [n]} X_i > \log n + 4x \right) = \mathbb{P}(N_n(-4x) \geq 1), \]
the Poisson approximation in Proposition 4.3 implies that for \( x \) large,
\[ \mathbb{P} \left( \max_{i \in [n]} X_i > \log n + 4x \right) \leq \frac{2(1 + o(1)) e^{-4x} \log n}{n} + (1 - \exp(-e^{-4x})) \]
\[ \leq (1 + o(1)) e^{-4x}. \] (4.85)

Further, by Proposition 4.9 for \( n \) large enough
\[ \mathbb{E}(R^{(1)}_n(x)) \leq C e^{-x/2}. \] (4.86)

It remains to analyse \( R^{(2)}_n(x) \). Arguing as in the proof of Proposition 4.9,
\[ \mathbb{E}(R^{(2)}_n(x)) \leq \mathbb{E}(N_n^2(-8x)) \mathbb{P}(d_w(1, 2) > \log n + 17x), \]
where \( d_w(1, 2) \) is the distance between vertices 1, 2 in \( \mathcal{S}_n = \{ K_n, (E_e)_{e \in \mathcal{E}_n} \} \). By Lemma 4.1,
\[ d_w(1, 2) = \frac{1}{N} \sum_{k=1}^{n} \frac{E_j}{k(n-k)}, \]
where \( N \) is uniform on \([n-1]\) independent of \( (E_j)_{j \in [n-1]} \) which are mean \( n \) exponential random variables. Thus, by Markov’s inequality, for any \( \alpha > 0 \),
\[ \mathbb{P}(d_w(1, 2) - \log n > 17x) \leq e^{-17\alpha} \sum_{j=1}^{n-1} \frac{1}{n-1} \exp \left( \alpha \left[ \log \frac{j}{n} - \log \left( 1 - \frac{j}{n} \right) \right] \right). \]
Letting $\beta = 1 - \epsilon$ with $\epsilon > 0$ small but independent of $x, n$, we finally get
\[
P(d_w(1, 2) - \log n > 17x) \leq (1 + o(1))e^{-17a\alpha}E\left(\left[\frac{U}{1 - U}\right]^{1 - \epsilon}\right),
\]
where $U \sim U[0, 1]$. We need to now bound $E(N_n^2(-8x))$. Write $N_n(-8x) = \sum_{i=1}^n Z_i$ where $Z_i = 1\{X_{(i)} \geq \log n + 8x\}$. By Proposition 4.3, $E(N_n(-8x)) \leq 2e^{8x}$. Further,
\[
\text{Var}(N_n(-8x)) \leq 2e^{8x} + n(n - 1)P(Z_1 = 1)\{P(Z_2 = 1|Z_1 = 1) - P(Z_2 = 1)\}.
\]
Given $Z_1 = 1$, the edge weights $(E_{2,i})_{i \neq 2}$ have the same distribution as
\[
(\log n - 8x + E_{2,1}, (E_{2,j})_{j \neq 1, 2}).
\]
Thus,
\[
P(Z_2 = 1|Z_1 = 1) = P\left(\min_{j \neq 2}E_{2,j} > \log n - 8x\right) = \exp\left(-\frac{n - 2}{n}(\log n - 8x)\right).
\]
Combining this, we get that $\text{Var}(N_n(8x)) \leq 4e^{8x}$ so that $E([N_n(-8x)]^2) \leq 16e^{16x}$. This results in
\[
E(R_w^2(x)) \leq (1 + o(1))16E\left(\left[\frac{U}{1 - U}\right]^{1 - \epsilon}\right)e^{-(1-17\epsilon)x}.
\]
Combining (4.85), (4.86) and (4.87) completes the proof of the asserted exponential tail bound in (4.83) and completes the proof of the upper bound.

\[\square\]

\textbf{Lower bound}. Let us now show that
\[
E\left[\left(\max_{i,j \in [n]} C_{ij} - 3\log n\right)^{2q}\right] = O(1).
\]
Recall that $V_1, V_2$ denote the vertices with the largest and second largest $X_{(i)}$ values. Further
\[
\max_{i,j \in [n]} C_{ij} - 3\log n \geq_{st} (X_{(V_1)} - \log n)_- + (X_{(V_2)} - \log n)_- + (d_w(1, 2) - \log n),
\]
where $d_w(1, 2)$ is independent of $X_{(V)}$ with the same distribution as the length of the optimal path between 1, 2 in $\mathcal{G}_n$ and $\geq_{st}$ denotes stochastic domination. By Hölder’s inequality,
\[
E\left[\left(\max_{i,j \in [n]} C_{ij} - 3\log n\right)^{2q}\right] \leq 3^{2q}\left(E([X_{(V_1)} - \log n]^{2q}) + E([X_{(V_2)} - \log n]^{2q}) + E([d_w(1, 2) - \log n]^{2q})\right).
\]
By [18, Proof of Theorem 3.3],
\[
E([d_w(1, 2) - \log n]^{2q}) = O(1).
\]
Further,
\[
E([X_{(V_1)} - \log n]^{2q}) \leq E([X_{(V_2)} - \log n]^{2q}).
\]
First recall the identity
\[
E(Y^{2q}) = 2q\int_0^\infty y^{2q-1}P(Y > y)\,dy,
\]

(4.89)
for any non-negative random variable $Y$. Second note that
\[
[X_{V_2} - \log n]_- \equiv 0, \quad \text{if } X_{V_2} \geq \log n. \tag{4.90}
\]
Suppose we show for some $0 < \varepsilon < 1$ small enough and fixed:
\[
P(\log n - X_{V_2} \geq x) \leq 2\exp^{-(1-\varepsilon)ex} + 2e^{2x\log n}\frac{n}{n}, \quad x \leq 2^{-1}(1 - \varepsilon)\log n. \tag{4.91}
\]
Then this bound implies that
\[
P(\log n - X_{V_2} \geq x) \leq C\frac{\log n}{n\varepsilon}, \quad x \in [2^{-1}(1 - \varepsilon)\log n, \log n]. \tag{4.92}
\]
Using these bounds on the tail probabilities in the moment identity (4.89) with $Y = [X_{V_2} - \log n]_-$ shows that
\[
\mathbb{E}([X_{V_2} - \log n]^{2q}) = O(1),
\]
which completes the proof. Thus we are now left with proving (4.91). However this follows from the Poisson approximation result in Proposition 4.3 since $P(\log n - X_{V_2} \geq x) = P(N_n(x) \leq 1)$. \square

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