Unique builders for classes of matrices *

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Abstract

Basic matrices are defined which provide unique building blocks for the class of normal matrices which include the classes of unitary and Hermitian matrices. Unique builders for quantum logic gates are hence derived as a quantum logic gates is represented by, or is said to be, a unitary matrix. An efficient algorithm for expressing an idempotent as a unique sum of rank 1 idempotents with increasing initial zeros is derived. This is used to derive a unique form for mixed matrices. A number of (further) applications are given: for example (i) $U$ is a symmetric unitary matrix if and only if it has the form $I - 2E$ for a symmetric idempotent $E$, (ii) a formula for the pseudo inverse in terms of basic matrices is derived. Examples for various uses are readily available.

1 Introduction

Basic matrices are defined and it is shown that any normal matrix is the product of basic commuting matrices and that the product is unique apart from the order. A matrix $B$ is normal when $BB^* = B^*B$ where $B^*$ denote the complex conjugate transposed of $B$. The class of normal matrices include the classes of unitary matrices ($UU^* = I$), and Hermitian, also called self-adjoint, matrices ($H^* = H$). These occur in many applications: for example quantum logic gates are represented by unitary matrices and their properties and applications depend ultimately on the structure of unitary matrices.

Each basic matrix itself is a product of minimal basic matrices with the same eigenvalue. A basic matrix is expressed in terms of a symmetric idempotent matrix; an idempotent $E$ occurring in the expression $B$ as a product of basic matrices has the property that $B$ acts on $E$ by $BE = \alpha E$ and the eigenvalue $\alpha$ occurs with multiplicity equal to the rank of $E$.

An efficient algorithm is given for expressing a rank $r$ idempotent matrix as a sum of rank 1 (pure) orthogonal idempotents with increasing initial zeros and such an expression is unique. From this a unique expression for a mixed matrix as a sum of rank 1 idempotents of this type is derived.

A quantum logic gate is represented by a unitary matrix and these gates are the basis for quantum information theory, see for example [3]. Indeed a quantum logic gate is often stated to be a unitary matrix itself. A quantum logic gate is thus a unique, apart from order, product of basic logic gates and these basic logic gates are building blocks for quantum logic gates in general.

Examples are readily available and applications are given throughout: It is shown that a unitary matrix $U$ is symmetric if and only if it has the form $U = I - 2E$ for a symmetric idempotent $E$. An easy formula for the roots and powers of the matrices follows directly from its expression as a product of basic matrices; powers and roots are given explicitly in terms of basic matrices. A formula for the pseudo inverse is immediate.

Expressions for well-known quantum gates (such as Pauli, Hadamard gates) as products of unique basic matrices are explicitly derived in section 4.

Comparisons may be made with the famous 1D factorization theorem of Belevich and Vaidyanathan which derives building blocks for 1D paraunitary matrices, [5] pp. 302-322. The basic matrices derived here for building normal matrices are influenced by methods in [2] for constructing generators/builders for multidimensional paraunitary matrices[1], and in particular for constructing paraunitary non-separable

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[1] A paraunitary matrix is a matrix $U(z)$ such that $U(z)U^*(z^{-1}) = I$ where $z = (z_1, z_2, \ldots, z_k)$ and $z^{-1} = (z_1^{-1}, z_2^{-1}, \ldots, z_k^{-1})$. 
(entangled) matrices.

A summary of the main results is given in Section 1.1. Further notation from Section 1.2 may be consulted as required.

1.1 Summary

Let $B$ be a normal matrix. Then $B$ is a product $B = \prod_{j=1}^{k} (I - E_j + \alpha_j E_j)$ the $\alpha_j$ are distinct and $S = \{E_1, E_2, \ldots, E_k\}$ is an orthogonal symmetric set of idempotents. Each $(I - E_j + \alpha E_j)$ is termed a basic matrix and the product is unique apart from the order of these commuting basic matrices.

$B$ acts on the idempotents by $BE_j = \alpha_j E_j$ and so $\alpha_j$ is an eigenvalue of $B$ occurring to multiplicity equal to rank $E_j$. When $S$ is not a complete set then $E = (I - \sum_{j=1}^{k} E_j)$ completes the set $S$ and 1 occurs as an eigenvalue of $B$ with multiplicity equal to $(n - \text{rank}(\sum_{j=1}^{k} E_j)) = n - (\sum_{j=1}^{k} \text{rank}(E_j))$ where $n$ is the size of the matrices.

A basic matrix $(I - E + \alpha E)$ is the product of minimal basic matrices $\prod_{j=1}^{k} (I - E_j + \alpha E_j)$ where the $E_j$ are mutually orthogonal, have rank 1 and each minimal basic has the same eigenvalue $\alpha$ as the original.

The idempotent matrix $E$ is of the form $E = \sum_{j=1}^{r} u_j u_j^*$ for mutually orthogonal unit column vectors $u_j$. Given an idempotent matrix $E$ of rank $r$ an efficient algorithm is given for expressing $E$ as a sum of such rank 1 idempotents with increasing initial zeros so that the expression obtained is unique. Mixed density matrices have a particularly useful form when viewed as products of basic matrices. A mixed density matrix is defined as a convex sum of pure density matrices though not in a unique way. By writing the mixed matrix as a product of basic matrices, uniqueness is obtained and this gives a unique perspective on mixed matrices. Further each basic matrix is a unique product of ordered rank 1 basics matrices giving a unique expression for a mixed density matrix. See Section 5.1.

Unitary matrices are normal matrices with eigenvalues of the form $e^{i\theta}$. Quantum logic gates are represented by unitary matrices; indeed a quantum logic gate is often defined to be a unitary matrix. The expression of a unitary matrix as a unique product, except for order, of basic matrices then expresses the original.

There are easily constructed interesting examples; those displayed are of small size but the constructions can be efficiently applied to large size matrices.

1.2 Additional notation

Necessary background on algebra may be found in many algebra and linear algebra books. Background on quantum information theory may be found in 3.

$R$ denotes a ring with identity $I_R$; the suffix $R$ may be omitted when a particular ring $R$ is understood. A mapping $^* : R \to R$ in which $r \mapsto r^*, (r \in R)$ is said to be an involution on $R$ if and only if (i) $r^{**} = (r^*)^* = r, \forall r \in R$, (ii) $(a + b)^* = a^* + b^*, \forall a, b \in R$, and (iii) $(ab)^* = b^*a^*, \forall a, b \in R$. Let $R$ be a ring with involution $^*$. An element $a \in R$ is said to be symmetric, with respect to $^*$, if $a^* = a$. An idempotent in $R$ is an element $E$ such that $E^2 = E$. $E, F$ are said to be orthogonal in $R$ if $EF = FE = 0_R$. The set $\{E_1, E_2, \ldots, E_k\}$ is said to be a complete set of orthogonal idempotents in $R$ if each element is an idempotent, the $E_i$ are mutually orthogonal and $E_1 + E_2 + \ldots + E_k = I_R$. The set is further said to be symmetric if each $E_i$ is symmetric (with respect to $^*$).
Here we work in the field of complex numbers $\mathbb{C}$ although many of the results work over other systems but are not included. For $a \in \mathbb{C}$, $a^*$ denotes the complex conjugate of $a$ and then $A^*$ denotes the complex conjugate transposed of $A$ for $A \in \mathbb{C}^{n \times m}$ or $A \in \mathbb{C}^n$. Now $I_n$ denotes the identity $n \times n$ matrix; the suffix $n$ will be omitted when the size is clear. An $n \times n$ matrix $B$ is said to be a normal matrix in $\mathbb{C}$ if and only if $BB^* = B^*B$. An $n \times n$ matrix is unitary if and only if $UU^* = I_n$. An $n \times n$ matrix is Hermitian (or self-adjoint) if and only if $H^* = H$. Thus unitary and Hermitian matrices are normal matrices. Unitary matrices by their definition are invertible but an Hermitian matrix may not be invertible.

Column vectors $u, v$ are orthogonal if $u^*v = 0$; $n \times n$ matrices are orthogonal if $AB^* = 0_{n \times n} = BA^*$; in all cases here orthogonality will refer to symmetric matrices. For a column $n \times 1$ vector $v \neq 0$, $vv^*$ is a symmetric $n \times n$ matrix and the matrix is necessarily of rank 1. When $v$ is a unit vector $(v^*v = 1)$ then $vv^*$ is an idempotent matrix as then $vv^*(vv^*)^* = vv^*vv^* = vv^*$. Suppose $v, w$ are orthogonal vectors. Then the matrices $vv^*$ and $ww^*$ are orthogonal matrices as $vv^*ww^* = 0_{n \times n}$ since $v^*w = 0$.

## 2 Basic Matrices

A matrix $A$ is normal if and only if there exists a unitary matrix $P$ such that $P^*AP = D$ where $D$ is a diagonal matrix. Normal matrices include Hermitian matrices ($H^* = H$) and unitary matrices ($UU^* = I$); Hermitian matrices include real symmetric matrices. Eigenvalues of a unitary matrix have the form $e^{i\theta}$ and an Hermitian matrix has real eigenvalues.

**Proposition 2.1** Let $A$ be an $n \times n$ matrix. There exists a unitary matrix $P$ such that $P^*AP = D$, where $D$ is diagonal, if and only if $A = \alpha_1 v_1 v_1^* + \alpha_2 v_2 v_2^* + \ldots + \alpha_n v_n v_n^*$ for an orthogonal set of unit vectors \( \{v_i| i = 1, 2, \ldots , n\} \).

This can be restated as follows:

**Proposition 2.2** Let $A$ be an $n \times n$ matrix. Then $A$ is a normal matrix if and only if $A = \alpha_1 v_1 v_1^* + \alpha_2 v_2 v_2^* + \ldots + \alpha_n v_n v_n^*$ for an orthogonal set of unit vectors \( \{v_i| i = 1, 2, \ldots , n\} \).

The $v_i$ consist of the columns of $P$ and the $\alpha_i$ are the diagonal entries of $D$ in Proposition 2.1 and are eigenvalues of $A$. Now \( \{v_1, v_2, \ldots , v_n\} \) is an orthonormal basis for $\mathbb{C}^n$. Define $E_i = v_i v_i^*$. Then $\{E_1, E_2, \ldots , E_n\}$ is an orthogonal symmetric set of idempotents. The set is also a complete set of idempotents for if not then $E = I - \sum_{i=1}^{n} E_i \neq 0$ and $E$ is orthogonal to each of the $E_i$; this would give a linearly independent set of $n + 1$ vectors in the $n$-dimensional space $\mathbb{C}^n$.

Now $\text{tr}A$ denotes the trace of the matrix $A$ which is the sum of its diagonal elements. One nice property of an idempotent matrix is that its rank is the same as its trace, see for example [1].

**Lemma 2.1** Let $E, F$ be orthogonal symmetric idempotent $n \times n$ matrices. Then

1. $E + F$ is a symmetric idempotent matrix.
2. If $E + F \neq I$, then $I - E - F$ is a symmetric idempotent matrix orthogonal to both $E$ and $F$.
3. $\text{rank}(E + F) = \text{rank} E + \text{rank} F$.
4. $\text{rank}(I - E - F) = n - \text{rank} E - \text{rank} F$.

**Proof:** The proofs of the first two items are straightforward. Proof of 3 It is known that $\text{rank} E = \text{tr}E$ for an idempotent matrix $E$, see for example [1]. Now $E + F$ is an idempotent and so $\text{rank}(E + F) = \text{tr}(E + F) = \text{tr}E + \text{tr}F = \text{rank} E + \text{rank} F$. The proof of 4 is similar. □

Lemma 2.1 may be generalised as follows:

**Lemma 2.2** Let $\{E_1, E_2, \ldots , E_s\}$ be a set of orthogonal symmetric idempotent matrices. Then

- $(E_1 + E_2 + \ldots + E_s)$ is an idempotent symmetric matrix.
• If $(E_1 + E_2 + \ldots + E_s) \neq I$ then $I - (E_1 + E_2 + \ldots + E_s)$ is a symmetric idempotent orthogonal to each $E_j$ for $j = 1, 2, \ldots, s$.

• $\text{rank}(E_1 + E_2 + \ldots + E_s) = \text{tr}(E_1 + E_2 + \ldots + E_s) = \text{tr}E_1 + \text{tr}E_2 + \ldots + \text{tr}E_s = \text{rank} E_1 + \text{rank} E_2 + \ldots + \text{rank} E_s$.

• $\text{rank}(I - (E_1 + E_2 + \ldots + E_s)) = n - \text{rank} E_1 - \text{rank} E_2 - \ldots - \text{rank} E_s$.

**Lemma 2.3** Let $\{E_1, E_2, \ldots, E_k\}$ be mutually orthogonal symmetric idempotent matrices. Then $(I - E_1 + \alpha_1 E_1)(I - E_2 + \alpha_2 E_2) \ldots (I - E_k + \alpha_k E_k) = \alpha_1 E_1 + \alpha_2 E_2 + \ldots + \alpha_k E_k + (I - E_1 - E_2 - \ldots - E_k)$.

The following follows directly from Lemma 2.3.

**Proposition 2.3** Let $A$ be an $n \times n$ matrix with $A = \alpha_1 E_1 + \alpha_2 E_2 + \ldots + \alpha_n E_n$ where $\{E_i|i = 1, 2, \ldots, n\}$ is a complete orthogonal set of symmetric idempotents. Then $A = (I - E_1 + \alpha_1 E_1)(I - E_2 + \alpha_2 E_2) \ldots (I - E_n + \alpha_n E_n)$.

A basic matrix is one of the form $(I - E + \alpha E)$ where $E$ is an idempotent. This has eigenvalue $\alpha$ occurring to multiplicity equal to rank $E$ and has eigenvalue 1 occurring to multiplicity equal to rank $(I - E) = n - \text{rank} E$, where $n$ is the size of $E$. When $\alpha \neq 0$, it has inverse $(I - E + \alpha^{-1}E)$.

**Lemma 2.4** Let $E, F$ be orthogonal symmetric idempotent matrices. Then $(I - E + \alpha E)(I - F + \alpha F) = I - (E + F) + \alpha(E + F)$

Lemma 2.4 enables basic matrices with orthogonal idempotents and the same eigenvalue to be collected together.

The following is a consequence of Lemma 2.4 and Proposition 2.3.

**Proposition 2.4** Let $A$ be an $n \times n$ matrix with $A = \alpha_1 E_1 + \alpha_2 E_2 + \ldots + \alpha_n E_n$ where $\{E_i|i = 1, 2, \ldots, n\}$ is a complete orthogonal set of symmetric idempotents. Moreover the eigenvalues of $A$ are $\{\alpha_i|i = 1, 2, \ldots, n\}$ and $\alpha_i$ occurs with multiplicity equal to rank $E_i$.

It is shown below, Proposition 2.4 is unique apart from the order of the commuting basic matrices. Note that it is required for this uniqueness that the $\alpha_i$ are distinct; the equal $\alpha_i$ have been gathered into one basic matrix by Lemma 2.4.

Each idempotent $E$ is a sum of rank 1 idempotents but not in a unique manner. Algorithm 3.1 below give a method of expressing an idempotent $E$ as the sum of rank 1 (pure) idempotents which have increasing initial zeros thus giving a unique expression for $E$ as a sum of such pure idempotents.

One or more of the $\alpha_j$ in Proposition 2.4 may be 0. Some of the $\alpha_j$ may be also 1 and then $(I - E_j + E_j) = I$ giving that $A = (I - E_1 + \alpha_1 E_1)(I - E_2 + \alpha_2 E_2) \ldots (I - E_k + \alpha_k E_k)$ where the $\alpha_j$ are distinct and $\neq 1$. Here then the eigenvalue 1 is in ‘disguise’ and if so, it occurs with multiplicity equal to $(n - \sum_{j=1}^{k} \text{rank} E_j)$.

Now $A$ acts on its idempotents by $AE_j = \alpha_j E_j$, and $AE = E$ when $E = (I - \sum_{j=1}^{k} E_j) \neq 0$.

In particular Proposition 2.4 can be applied to unitary matrices. The eigenvalues of a unitary matrix are of the form $e^{i\theta}$.

**Proposition 2.5** Let $U$ be an $n \times n$ unitary matrix. Then $U = (I - E_1 + e^{i\theta_1} E_1)(I - E_2 + e^{i\theta_2} E_2) \ldots (I - E_k + e^{i\theta_k} E_k)$, the $\alpha_j$ are distinct and $\{E_1, E_2, \ldots, E_k\}$ is an orthogonal set of idempotents. Moreover the eigenvalues of $U$ are $\{e^{i\theta_i}, i = 1, 2, \ldots, k\}$ and $e^{i\theta_i}$ occurs with multiplicity rank $E_i$. Now eigenvalue 1 occurs with multiplicity $n - (\sum_{j=1}^{k} \text{rank} E_j)$ which may be 0. Moreover a matrix of the form $(I - E_1 + e^{i\theta_1} E_1)(I - E_2 + e^{i\theta_2} E_2) \ldots (I - E_k + e^{i\theta_k} E_k)$ is a unitary matrix for any orthogonal set of symmetric idempotents $\{E_1, E_2, \ldots, E_k\}$.

A basic unitary matrix is one of the form $(I - E + e^{i\theta} E)$. 

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In Proposition 2.8, A may be singular in which case the eigenvalue 0 appears to a certain multiplicity. In this case $A = (I - E_1 + \alpha_1 E_1)(I - E_2 + \alpha_2 E_2) \cdots (I - E_k + \alpha_k E_k)(I - E)$, where the $\alpha_i$ are distinct and $\neq 0$ and $\{E_1, E_2, \ldots, E_k, E\}$ is an orthogonal set of idempotents. (The eigenvalue 1 is hidden when $\{E_1, E_2, \ldots, E_k, E\}$ is not complete.) It is easy to check that $A^* = \prod_{j=1}^k (I - E_j + \alpha_j^* E_j)(I - E)$ is the pseudo inverse of A.

**Proposition 2.6** Let $A = (I - E_1 + \alpha_1 E_1)(I - E_2 + \alpha_2 E_2) \cdots (I - E_k + \alpha_k E_k)(I - E)$, where the $\alpha_i$ are distinct and $\neq 0$ and $\{E_1, E_2, \ldots, E_k, E\}$ is an orthogonal set of idempotents. Then $A^* = \prod_{j=1}^k (I - E_j + \alpha_j^* E_j)(I - E)$ is the pseudo inverse of A.

**Proposition 2.7** Let $A$ be a normal matrix. Then $A^2 = I$ if and only if $A = (I - 2E)$ for a symmetric idempotent E.

**Proof:** Suppose $A^2 = I$. Then the eigenvalues of A are $\{1, 1\}$. Thus by Proposition 2.4 $A = (I - E - E)(I - F + F) = (I - 2E)$ for symmetric orthogonal idempotents $E,F$. If $A = (I - 2E)$ then clearly $A^2 = I$ and A can be diagonalised by a unitary matrix.

The following is a corollary:

**Proposition 2.8** $U$ is a symmetric $(U^* = U)$ unitary matrix if and only if $U = I - 2E$ for an idempotent $E$.

Notice in Proposition 2.7 and in Proposition 2.8 the eigenvalue $\{-1\}$ occurs with multiplicity equal to rank $E$ and the eigenvalue $\{1\}$ occurs with multiplicity equal to rank $F = n - \text{rank} E$.

**Corollary 2.1** Let $HH = nI$ and $H = H^*$ for an $n \times n$ matrix $H$. Then $H = \sqrt{n}(I - 2E)$ for a symmetric idempotent $E$. In particular this applies to a symmetric Hadamard matrix.

### 2.1 Examples

1. $U = \left( \begin{array}{cc} 0 & 1 \\ 1 & 0 \end{array} \right)$. This has eigenvalues $\{1, -1\}$ and corresponding eigenvectors $(1, 1)^T, (1, -1)^T$. Then $U = E_1 - E_2 = E_1 + e^{i\theta} E_2 = (I - E_1 + E_1)(I - E_2 + e^{i\theta} E_2) = I - E_2 + e^{i\theta} E_2$ where $E_1 = \frac{1}{2} \left( \begin{array}{c} 1 \\ 1 \end{array} \right), E_2 = \frac{1}{2} \left( \begin{array}{c} 1 \\ -1 \end{array} \right)$. Using this formula, roots of $U$ may be obtained directly – see Section 3.2.

2. A commonly referenced matrix is the following real orthogonal/unitary matrix $U = \left( \begin{array}{cc} \cos \theta & \sin \theta \\ -\sin \theta & \cos \theta \end{array} \right)$. This has eigenvalues $\{e^{i\theta}, e^{-i\theta}\}$. Then $U = e^{i\theta} E_1 + e^{-i\theta} E_2 = \frac{1}{2} e^{i\theta} \left( \begin{array}{c} 1 \\ -1 \end{array} \right) + \frac{1}{2} e^{-i\theta} \left( \begin{array}{c} 1 \\ 1 \end{array} \right)$, which may be checked independently. The $E_i$ are symmetric idempotents. It is seen that $U = (I - E_1 + e^{i\theta} E_1)(I - E_2 + e^{-i\theta} E_2)$ which gives $U$ as a product of basic unitary matrices, which are defined in Section 2.2 below.

### 2.2 Uniqueness

**Definition** A basic matrix is one of the form $B = (I - E + \alpha E)$ for a symmetric idempotent matrix $E$.

The idempotent $E$ of $B$ is said to be the idempotent involved in $B$ and $\alpha$ is the eigenvalue involved in $B$.

**Definition** Say a basic matrix $I - E + \alpha E$ is a simple basic unitary matrix if the idempotent $E$ has rank 1.

**Proposition 2.9** A symmetric idempotent $E$ has rank 1 if and only if $E = uu^*$ for a unit column vector $u$.

**Proof:** If $E$ has the form $uu^*$ for a unit column vector $u$, then clearly $E$ is a symmetric idempotent of rank 1. On the other hand if $E$ is a symmetric idempotent of rank 1 then $E$ is of the form $uu^*$ for a unit column vector $u$. This is shown in for example [2] Proposition 4.5.

A minimal basic is of the form $(I - E + \alpha E)$ for $E = uu^*$ where $u$ is a unit column vector and has rank 1 and $\alpha \neq 0$. 

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Note Let \( I - E + \alpha E \) and \( I - E + \beta E \) be basic matrices with the same idempotent. Then \((I - E + \alpha E)(I - E + \beta E) = I - E + \alpha\beta E\). A product of a minimal basic with another minimal basic with the same idempotent is another minimal basic. The minimality is expressed in terms of the rank of the idempotent involved.

**Proposition 2.10** A symmetric idempotent \( E \) has rank \( k \) if and only if \( E = u_1u_1^* + u_2u_2^* + \ldots + u_ku_k^* \) for unit column mutually orthogonal vectors \( u_i \).

**Proof:** When \( k = 1 \) the result follows from Proposition 2.9 If the first column of \( E \) is zero then the first row of \( E \) is zero and the result follows by induction. Suppose then the first column \( w \) of \( E \) is non-zero and define \( u = w/|w| \) which is a unit column vector. Then \( F = uu^* \) is an idempotent of rank 1. Let \( E_1 = E - F \). Then \( E_1 = E - F \) is an idempotent symmetric matrix orthogonal to \( E \) which has first column, and hence first row, consisting of zeros. Let \( A \) be the \((n-1) \times (n-1)\) matrix with first column and first row of \( E_1 \) omitted. Since \( F, E_1 \) are orthogonal it follows that rank \( F + \text{rank} E_1 = \text{rank} (F + E_1) = \text{rank} E \) and so \( E - F \) has rank \((k-1)\). Thus \( A \) is rank \((k-1)\) symmetric idempotent matrix. The result then follows by induction.

Thus a basic unitary matrix is of the form \((I - E + \alpha E)\) where \( E = u_1u_1^* + u_2u_2^* + \ldots + u_ku_k^* \) for unit column mutually orthogonal vectors \( u_i \). When \( k = 1 \) the basic unitary matrix is what is termed a simple basic unitary matrix and its idempotent is of the form \( uu^* \).

**Proposition 2.11** Suppose the space \( W \) is generated by the unit orthogonal columns vectors \( \{w_1, w_2, \ldots, w_k\} \) and by the unit orthogonal column vectors \( \{v_1, v_2, \ldots, v_k\} \). Then

\[
w_1w_1^* + w_2w_2^* + \ldots + w_kw_k^* = v_1v_1^* + \ldots + v_kv_k^*.
\]

**Proof:** Now

\[
w_1 = \alpha_{11}v_1 + \alpha_{12}v_2 + \ldots + \alpha_{1k}v_k \\
w_2 = \alpha_{21}v_1 + \alpha_{22}v_2 + \ldots + \alpha_{2k}v_k \\
\vdots \\
w_k = \alpha_{k1}v_1 + \alpha_{k2}v_2 + \ldots + \alpha_{kk}v_k
\]

for some \( \alpha_{ij} \in \mathbb{C} \).

In other words

\[
(w_1, w_2, \ldots, w_k) = (v_1, v_2, \ldots, v_k)
\]

\[
\begin{pmatrix}
\alpha_{11} & \alpha_{12} & \ldots & \alpha_{1k} \\
\alpha_{21} & \alpha_{22} & \ldots & \alpha_{2k} \\
\vdots & \vdots & \ddots & \vdots \\
\alpha_{k1} & \alpha_{k2} & \ldots & \alpha_{kk}
\end{pmatrix}
\]

Denote the matrix on the right (involving \( \alpha_{ij} \)) by \( U \).

Then

\[
(w_1, w_2, \ldots, w_k) \begin{pmatrix}
w_1^* \\
w_2^* \\
\vdots \\
w_k^*
\end{pmatrix} = (v_1, v_2, \ldots, v_k)UU^*
\]

Therefore

\[
(w_1, w_2, \ldots, w_k) \begin{pmatrix}
w_1^* \\
w_2^* \\
\vdots \\
w_k^*
\end{pmatrix} = (v_1, v_2, \ldots, v_k)UU^*
\]

In conclusion
Then by (1),

\[
I_k = \begin{pmatrix} w_1^* \\ w_2^* \\ \vdots \\ w_k^* \end{pmatrix} (w_1, w_2, \ldots, w_k) = U^* \begin{pmatrix} v_1^* \\ v_2^* \\ \vdots \\ v_k^* \end{pmatrix} (v_1, v_2, \ldots, v_k) U = U^* U
\]

Thus \(U^* U = I_k\) and hence \(UU^* = I_k\). Therefore by (2)

\[
(w_1, w_2, \ldots, w_k) \begin{pmatrix} w_1^* \\ w_2^* \\ \vdots \\ w_k^* \end{pmatrix} = (v_1, v_2, \ldots, v_k) \begin{pmatrix} v_1^* \\ v_2^* \\ \vdots \\ v_k^* \end{pmatrix}
\]

and hence \(w_1 w_1^* + w_2 w_2^* + \ldots + w_k w_k^* = v_1 v_1^* + \ldots + v_k v_k^*\).

\[\Box\]

Suppose now \(B = (1 - E_1 + \alpha_1 E)(1 - E_2 + \alpha_2 E_2) \ldots (1 - E_k + \alpha_k E_k)\) is a product of basic \(n \times n\) matrices. Implicit in this is that if \(r = \text{rank}(E_1 + \text{rank}(E_2 + \ldots + \text{rank}(E_k) < n\) then \(F = (I - E_1 - E_2 - \ldots - E_k)\) is a non-zero idempotent with eigenvalue 1 of multiplicity \((n - r)\). By Lemma 2.4 a product basic matrices using orthogonal idempotents with the same eigenvalues can be collected together into a basic matrix.

Suppose \(B = (I - F_1 + \alpha_1 F_2)(I - F_2 + \alpha_2 F_2) \ldots (I - F_k + \alpha_k F_k)\) where the \(F_j\) are mutually orthogonal idempotents and the \(\alpha_j\) are all different. We now show that except for the order such an expression is unique. Each \(F_j\) in the product has the form \(w_k w_k^* + w_2 w_2^* + \ldots + w_k w_k^*\) where \(k\) is the rank of \(F_j\) and is also the multiplicity of the corresponding eigenvalue \(\alpha_j\).

**Proposition 2.11** Suppose \(B = (I - E_1 + \alpha_1 E_1)(I - E_2 + \alpha_2 E_2)(\ldots)(I - E_k + \alpha_k E_k) = (I - F_1 + \beta_1 F_1)(I - F_2 + \beta_2 F_2)(\ldots)(I - F_k + \beta_k F_k)\) where the \(E_j\) are orthogonal idempotents and the \(\alpha_j\) are distinct, and where the \(E_j\) are orthogonal idempotents and the \(\beta_j\) are distinct. Then \(s = k\) and by reordering \(F_1 = E_1, \alpha_1 = \beta_1\).

**Proof:** Now \(B\) has eigenvalues \(\alpha_j\) occurring with multiplicity rank \(E_j\) and eigenmatrix \(E_j\) and looking at it another way has eigenvalues \(\beta_j\) occurring with multiplicity rank \(F_k\). Thus \(\alpha_1\) must equal \(\beta_1\) for some \(k\). We may assume this \(k = 1\) by reordering. Then rank \(E_1 = \text{rank} F_1\). Now \(E_1 = v_1 v_1^* + v_2 v_2^* + \ldots + v_t v_t^*\) for orthogonal unit \(v_j\), \(t = \text{rank} E_1 = \text{rank} F_1\) and \(F_1 = w_1 w_1^* + w_2 w_2^* + \ldots + w_t w_t^*\) for orthogonal unit vectors \(w_j\). Now by Proposition 2.11 \(E_1 = F_1\). Similarly by reordering if necessary \(E_j = F_j, \alpha_j = \beta_j\) in general for \(j = 1, 2, \ldots, k\) and of necessity \(s = k\). \[\Box\]

## 3 Pure and mixed idempotents

In order to obtain a normal matrix as a unique product of basic matrices it is necessary to collect the basics with the same eigenvalue, see Propositions 2.4 and 2.12. An idempotent is not a unique sum of idempotents of rank 1. An idempotent of rank 1 is often described as a pure idempotent by analogy with the mixed matrix in quantum theory, see Section 5.1 below.

It is now shown that an idempotent matrix \(E\) of rank \(r\) may be written as the sum of special idempotents of rank 1 in a unique way and an efficient algorithm is given.\[\footnote{Induction is more complicated because we cannot assume \((I - E_1 + \alpha_1 E)\) is invertible.} \]
Let $E$ be an $n \times n$ idempotent matrix of rank $r$. Then $(I - E)$ is an idempotent matrix of rank $(n - r)$. Also $E$ has eigenvalues $\{1, 0\}$. Now $E.E = E$ and so the multiplicity of the eigenvalue 1 is $\geq r$ since $E$ has rank $r$. Also $E.(I - E) = 0 = 0(I - E)$ and so the multiplicity of the eigenvalue 0 is $\geq (n - r)$ as $(I - E)$ has rank $(n - r)$. Hence the multiplicity of the eigenvalue 1 of $E$ is exactly $r$ and the multiplicity of the eigenvalue 0 of $E$ is exactly $(n - r)$.

The following gives $E$ as a sum of rank 1 idempotents from the column space of $E$.

**Proposition 3.1** Let $\{e_1, e_2, \ldots, e_r\}$ be an orthonormal basis for the column space of $E$. Define $E_i = e_i e_i^*$ for $i = 1, 2, \ldots, r$. Then $E = E_1 + E_2 + \ldots + E_r$.

**Proof:** The space of eigenvectors of $E$ corresponding to the eigenvalue $\lambda = 1$ is the solution space of $(E - I)\mathbf{v} = 0$ which is the solution space of $(I - E)\mathbf{v} = 0$ and this has dimension $r$. Now $(I - E)E = 0$ and so the columns of $E$ form a basis for this solution space. Let $\{v_1, v_2, \ldots, v_r\}$ be an orthonormal basis for this solution space which is an orthonormal basis for the columns of $E$.

The space of eigenvectors of $E$ corresponding to the eigenvalue $\lambda = 0$ is the solution space of $(E - 0I)\mathbf{v} = 0$ which is the solution space of $E\mathbf{v} = 0$ and this has dimension $(n - r)$. Now $E(I - E) = 0$ and so the columns of $I - E$ form a basis for this solution space. Let $\{v_{r+1}, v_{r+2}, \ldots, v_n\}$ be an orthonormal basis for this solution space which is an orthonormal basis for the column space of $(I - E)$.

Now vectors corresponding to distinct eigenvalues of the symmetric matrix $E$ are orthogonal and so $\{v_1, v_2, \ldots, v_n\}$ is an orthonormal basis for $C_n$. Let $P = (v_1, v_2, \ldots, v_n)$. Then $P$ is a unitary matrix and $P^* E P = \text{diag}(1, 1, \ldots, 1, 0, 0, \ldots, 0)$ where the 1 occurs $r$ times and the 0 occurs $(n - r)$ times on the diagonal. Then $E = v_1 v_1^* + v_2 v_2^* + \ldots + v_r v_r^*$ as required. □

This expression for $E$ is not unique as any orthonormal basis for the column space of $E$ may be used.

An efficient algorithm for finding $E$ as a sum of idempotents of rank 1 is then formulated as follows:

**Algorithm 3.1**

1. Let $v$ be the first non-zero column of $E$ which occurs at column $r \geq 1$ so that columns $\{1, \ldots, r - 1\}$ consist of zeros.

2. Define $u = \frac{v}{||v||}$ so $u$ is a unit column vector.

3. Define $E_1 = uu^*$. Then $E_1$ is rank 1 symmetric idempotent.

4. Let $F = E - E_1$. If $F = 0$ then $E = E_1$ and $E$ has rank 1 and we are finished. If $F \neq 0$ then $F$ is a nonzero idempotent with columns $1, 2, \ldots, r$ all zero vectors. Now let $E$ be replaced by $F$ and go back to item 1.

To apply the algorithm it is not necessary to know the rank of the idempotent beforehand and this comes out at the end. The algorithm is efficient and at each step just involves forming an idempotent from the first non-zero column of a matrix and subtracting this from the matrix.

**Example 1** $E = \frac{1}{3} \begin{pmatrix} 2 & -1 & 1 \\ -1 & 2 & 1 \\ 1 & 1 & 2 \end{pmatrix}$. Then $E$ is a rank 2 idempotent. The first column of $E$ is non-zero so start there and let

$E_1 = \frac{1}{3} \begin{pmatrix} 2 \\ -1 \\ 1 \end{pmatrix} (2, -1, 1) = \frac{1}{3} \begin{pmatrix} 4 & -2 & 2 \\ -2 & 1 & -1 \\ 2 & -1 & 1 \end{pmatrix}$

Then $E - E_1 = \frac{1}{3} \begin{pmatrix} 0 & 0 & 0 \\ 0 & 1 & 1 \\ 0 & 1 & 1 \end{pmatrix} = E_2$. Now $E_2$ is an idempotent of rank 1 and $E = E_1 + E_2$. 

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Example 2  Let \( E = \frac{1}{4} \begin{pmatrix} 3 & i & 1 & -i \\ -i & 3 & i & 1 \\ 1 & -i & 3 & i \\ i & 1 & -i & 3 \end{pmatrix} \).

Apply the algorithm and get in turn \( E_1, E_2, E_3 \) idempotents of rank 1 and \( E = E_1 + E_2 + E_3 \) where
\[
E_1 = \frac{1}{12} \begin{pmatrix} 9 & 3i & 3 & -3i \\ -3i & 1 & -i & 1 \\ 3 & i & 1 & -i \\ 3i & -1 & i & 1 \end{pmatrix}, \quad E_2 = \frac{1}{6} \begin{pmatrix} 3 & 0 & 0 & 0 \\ 0 & 4 & 2i & 2 \\ 0 & -2i & 1 & -i \\ 0 & 2 & i & 1 \end{pmatrix}, \quad E_3 = \frac{1}{8} \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & i \\ 0 & 0 & -i & 1 \end{pmatrix}.
\]

\( E_1 \) is formed using the first (non-zero) column of \( E \) and notice that the number of initial zeros increase from \( E_1 \) to \( E_2 \) to \( E_3 \).

For a mixed matrix \( A = \sum_i p_i v_i v_i^* \) with \( 0 < p \leq 1, \sum p_i = 1 \), as in Section 3.1 Now as shown \( A = \sum_i p_i E_i \) where the \( p_i \) are distinct, each \( E_i \) is a basic idempotent and this expression for \( A \) is unique apart from the order of the \( E_i \). Each \( E_i \) is a sum of of rank 1 idempotents in a non-unique manner and such a sum may be obtained from Proposition 3.1 in theory and by Algorithm 3.1 in an efficient algorithm.

Let \( u = (\alpha_1, \alpha_2, \ldots, \alpha_n)^T \) be a column vector. Say \( sI(u) = 0 \) if \( \alpha_1 \neq 0 \) (if initial entry of \( u \) is non-zero) and say \( sI(u) = r \) if \( \alpha_1 = \alpha_2 = \ldots = \alpha_r \) but \( \alpha_{r+1} \neq 0 \) (if \( \alpha_{r+1} \) is the first non-zero entry of \( u \)). Let \( E = uu^* \) then define \( sI(E) = sI(u) \). Thus \( E \) has first \( sI(E) \) columns and first \( sI(E) \) rows consisting entirely of zeros. The \( sI \) measures the number of initial zeros of \( u \) and the number of initial rows and columns of \( E \).

Theorem 3.1  Let \( E \) be an idempotent. Then \( E \) is uniquely a sum \( E = E_1 + E_2 + \ldots + E_r \) where each \( E_i \) is a rank 1 idempotent and \( sI(E_1) < sI(E_2) < \ldots < sI(E_r) \).

Proof:  Algorithm 3.1 shows that has such an expression for \( E \) exists. The process involves constructing first an idempotent \( E_1 \) such that \( E - E_1 \) is an idempotent but also that \( sI(E - E_1) < sI(E) \). Now work with \( (E - E_1) \) and proceed until the zero matrix is obtained.

Suppose now \( E = E_1 + E_2 + \ldots + E_r = F_1 + F_2 + \ldots + F_s \) where \( sI(E_1) < sI(E_2) < \ldots < sI(E_r) \) and \( sI(F_1) < sI(F_2) < \ldots < sI(F_s) \) where the \( E_i \) and \( F_j \) are rank 1 idempotents. Then \( r = s \) as this is the rank of \( E \). Now look at the first non-zero column of \( E_1 \) which is the first non-zero column of \( E \). From this get \( E_1 = F_1 \) and then by induction that \( E_i = F_i \) for \( i = 1, 2, \ldots, r \).

These results are applied below in Section 3.1 to derive a unique expression for a mixed matrix as the sum of rank 1 matrices of the form given in Theorem 3.1.

3.1 Mixed matrix

Suppose \( A = \sum_{i=1}^n p_i u_i u_i^* \) with \( p_i \neq 0 \) where \( \{u_1, u_2, \ldots, u_k\} \) are mutually orthogonal unit (column) vectors in \( C_n \). Complete \( \{u_1, u_2, \ldots, u_k\} \) to an orthogonal unit basis \( \{u_1, u_2, \ldots, u_n\} \) for \( C_n \). Let \( E_i = v_i v_i^* \) and \( p_j = 0 \) for \( j > k \). Now \( A = \sum_{i=1}^n p_i u_i u_i^* = \sum_{i=1}^n p_i E_i = \prod_{j=1}^k (I - E_j + p_j E_j) (I - E) \) where \( \{E_1, E_2, \ldots, E_k, E\} \) is a complete orthogonal set of idempotents. Basic matrices with the same \( p_i \) may be collected together. Note that \( E = I - \sum_{j=1}^k E_j \) as \( \{E_1, E_2, \ldots, E_k, E\} \) is a complete orthogonal set of idempotents.

It is easy to check that \( A^+ = \prod_{j=1}^k (I - E_j + p_j^{-1} E_j) (I - E) \) is the pseudo inverse of \( A \) which is useful for certain calculations.

A particular case of this is where \( A \) is a mixed density matrix \( A = \sum_{i=1}^k p_i u_i u_i^* \) where \( \{u_1, u_2, \ldots, u_k\} \) are orthogonal unit vectors and \( \sum p_i = 1, 0 < p_i \leq 1 \). It is known that this expression for \( A \) is not unique.

But \( A = \sum_{i=1}^k p_i E_i \) where the \( E_i = u_i u_i^* \) are rank 1 idempotents. Then \( A = \prod_{i=1}^k (I - E_i + p_i E_i) (I - F) \) where \( F = I - \sum_j^k E_j \). Now by Lemma 2.12 the basic products with the same \( p_i \) can be gathered together to get \( A = \prod_{i=1}^k (I - F_i + \tilde{p}_i F_i) (I - F) \) where \( \{F_1, F_2, \ldots, F_s, F\} \) is a complete orthogonal set of idempotents with \( F = I - \sum_j^s F_j \), the \( \tilde{p}_i \) are unique and \( \sum_{j=1}^s \tilde{p}_j = 1, 0 < \tilde{p}_j \leq 1 \). This expression for \( A \) is unique by Proposition 2.12. This means that \( A = \sum_{i=1}^k \tilde{p}_i F_i \) where the \( \{F_j\} \) are orthogonal
idempotents and the \( \tilde{p}_j \) is now a unique expression for \( A \). Each \( F_j \) is a sum \( \sum_{i=1}^{s_j} u_i u_i^* \) of orthogonal rank 1 idempotents and \( F_j \) is of rank \( s_j \).

Note that \( AF_j = \tilde{p}_j F_j \) and \( \tilde{p}_j \) occurs with multiplicity equal to rank \( F_j \).

Using Theorem \( \ref{thm:subspace} \) the following unique expression is obtained.

**Theorem 3.2** A density matrix \( A \) is uniquely of the form \( A = \sum_{j} \lambda_j (E_{j,1} + E_{j,2} + \ldots + E_{j,s_j}) \) with distinct \( \lambda_j \) and such \( \text{st}(E_{j,1}) < \text{st}(E_{j,2}) < \ldots < \text{st}(E_{j,s_j}) \) for each \( j \).

### 3.2 Roots

A nice formula for roots of normal matrices may now be obtained directly from their expression as a product of basic matrices. Roots of unitary matrices/quantum gates will then follow.

Powers of normal matrices are easy to obtain:

**Proposition 3.2** Let the normal matrix \( B \) be expressed as a product of basic matrices by \( B = \prod_{j=1}^{k} (I - E_j + \alpha_j E_j) \). Then \( B^n = \prod_{j=1}^{k} (I - E_j + \alpha^n E_j) \).

**Proof:** The proof follows directly by induction and noting that the \( E_j \) commute. \( \square \)

Let \( U = \prod_{j=1}^{k} (I - E_j + \alpha_j E_j) \) be the expression of an \( n \)th root of \( I_n \) as a product of basic matrices. The \( E_j \) are mutually orthogonal idempotents and the \( \alpha_j \neq 0 \).

Then \( U^n = \prod_{j=1}^{k} (I - E_j + \alpha^n E_j) = I_n \). Hence \( \alpha^n = 1 \) and so \( \alpha \) is an \( n \)th root of unity. Hence \( \alpha_j = e^{i\theta_j} \) where \( n\theta_j = 2\pi r_j \) for \( r_j \in \mathbb{Z}, 1 \leq r_j \leq n - 1 \). Thus \( \theta_j = \frac{2\pi r_j}{n} \) and so

\[
U = \prod_{j=1}^{k} (I - E_j + e^{\frac{2\pi r_j}{n} E_j})
\]

This gives normal matrices which the \( n \)th roots of \( I_n \). The \( E_j \) can be any set of orthogonal idempotents and the \( r_j \) can be any integers between 1 and \( n - 1 \).

**Proposition 3.3** The \( n \)th roots of \( I_n \) which are normal consist of matrices of the following type:

\[
U = \prod_{j=1}^{k} (I - E_j + e^{i\frac{2\pi r_j}{n} E_j})
\]

where \( \{E_1, E_2, \ldots, E_k\} \) is an orthogonal set of idempotents in \( \mathcal{C}_{n \times n} \) and \( r_j \) are integers satisfying \( 1 \leq r_j \leq n - 1 \).

Suppose \( V = \{E_1, E_2, \ldots, E_k\} \) is a set of orthogonal symmetric idempotent matrices in \( \mathcal{C}_{n \times n} \). If \( V \) is not complete then there exists \( S = \{E_{k+1}, E_{k+2}, \ldots, E_s\} \) such that \( \{E_1, E_2, \ldots, E_s\} \) is a complete orthogonal set of idempotents. The set \( S \) is not unique but we shall refer to an \( S \) as a complementary set of \( V \) and say that \( S \) completes \( V \). In particular \( \{I - \sum_{j=1}^{k} E_j\} \) is the minimal set which completes \( V \).

Suppose \( U \) is the \( n \)th root of a basic matrix \( W \) and \( W = \prod_{j=1}^{k} (I - F_j + \beta_j F_j) \) is the expression for \( W \) as a product of basic matrices where the \( \beta_j \) are all different. Let \( U = \prod_{j=1}^{k} (I - E_j + \alpha_j E_j) \) be the expression for \( U \) as a product of basic matrices. Some of the \( (I - E_j + \alpha_j E_j) \) may be \( n \)th roots of \( I_n \) and we separate these out so that \( U = \prod_{j=1}^{k} (I - E_j + \alpha_j E_j)P \) where \( P \) is an \( n \)th root of \( I_n \) and commutes with each \( E_j \); reordering the \( E_j \) may be necessary. Now \( P \) is a product of basic matrices involving only idempotents in a complementary set of \( \{E_1, E_2, \ldots, E_s\} \).

Then \( U^n = \prod_{j=1}^{k} (I - E_j + \alpha^n E_j)P^n = \prod_{j=1}^{k} (I - E_j + \alpha^n E_j) \). Now compare this with \( W \). By uniqueness and possible reordering get that \( E_j = F_j \) and \( \alpha_j^n = \beta_j \). Thus if \( \alpha_j = r_j e^{i\gamma_j} \) and \( \beta_j = s_j e^{i\gamma_j} \) then \( r_j = s_j \) and \( n\beta_j = \gamma_j + 2\pi r_j \). Now \( n\alpha_j = 2\pi r_j + s_j \) for \( 0 < s_j < 2\pi \) and so \( e^{i\gamma_j} = e^{i\gamma_j} \).

However a normal matrix has many \( n \)th roots and we now proceed to directly find all these.
Suppose $U = \prod_{j=1}^{k}(I - E_j + \alpha_j E_j)$. Now $\sum E_j$ may not be $I$ and in this case define $F = (I - \sum E_j)$. Then \{\$E_1, E_2, \ldots, E_n, F\$\} is a complete orthogonal set of idempotents. Now $I = I + e^{i\theta}F$ and we need to attach this to $U$: $U = \prod_{j=1}^{k}(I - E_j + \alpha_j E_j)(I - F + e^{i\theta}F)$. Suppose $n^{th}$ roots of $U$ are required. Apply the usual way of obtaining an $n^{th}$ root of unity in $C$ and write: $U = \prod_{j=1}^{k}(I - E_j + \alpha_j e^{i2rj\pi/k} E_j)(I - F + e^{i2r\pi/k}F)$ for positive integers $r_j, r$ from 0 to $n - 1$. Then $U^{n} = \{\prod_{j=1}^{k}(I - E_j + \alpha_j e^{i12rj\pi/k} E_j)(I - F + e^{i2r\pi/k}F)\}$ and this is for any $r_j, r$ between 0 and $n - 1$. Here $\alpha^n$ refers to real $n^{th}$ of the the modulus of $\alpha$. This gives all the $n^{th}$ roots of $U$ and they are expressed in a unique form as a product of basic matrices. If the expression for $U$ has $k + 1$ basic matrices, including $F \neq 0$, then get $n^{k+1}$ different $n^{th}$ roots of $U$.

### 3.2.1 Examples:

Consider the examples 2.1.

1. Example 2 uses the matrix $U = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ with $E_1 = \frac{1}{2}(\begin{smallmatrix} 1 & 1 \\ 1 & -1 \end{smallmatrix})$ and $U = I - E_2 + e^{i\pi}E_2$. The square roots of $U$ are then $\{I - E_2 + e^{i\frac{\pi}{2}}E_2, I - E_2 + e^{i\frac{3\pi}{2}}E_2\}$ plus a product of these using any square roots of $I$ obtained from a complementary set of \{\$E_2\$\}. The only complementary set of \{\$E_2\$\} is \{\$E_1\$\}. A basic matrix involving $E_1$, besides the identity, which is a square root of $I$ is $I - E_1 + e^{i\pi}E_1$. Thus the square roots of $U$ are $\{I - E_2 + e^{i\frac{\pi}{2}}E_2)(I - E_1 + e^{i\pi}E_1), (I - E_2 + e^{i\frac{3\pi}{2}}E_2)(I - E_2 + e^{i\pi}E_2)\}$ where $k, s$ can be 0 or 1.

2. Example 2 uses the orthogonal matrix $U = \begin{pmatrix} \cos\theta & \sin\theta \\ -\sin\theta & \cos\theta \end{pmatrix}$. Then $U = (I - E_1 + e^{i\theta}E_1)(I - E_2 + e^{-i\theta}E_2)$ with \{\$E_1 = \frac{1}{2}(\begin{smallmatrix} 1 & 0 \\ 0 & 1 \end{smallmatrix})$, $E_2 = \frac{1}{2}(\begin{smallmatrix} 0 & 1 \\ 1 & 0 \end{smallmatrix})\}$. Here \{\$E_1, E_2\$\} is a complete orthogonal set of idempotents so there is no complementary set to consider.

The third roots of $U$ are $\{(I - E_1 + e^{i\frac{2\pi}{3}} E_1)(I - E_2 + e^{-i\frac{2\pi}{3}} E_2)\}$ where $r, s$ can have values from \{0, 1, 2\}.

### 4 Logic gates

Expressions as products of basic matrices are found for common quantum logic gates.

- **Hadamard Gate** The unitary matrix of the the Hadamard Gate is $H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$. Then $H^2 = I$ and $H$ has eigenvalues \{1, -1\}. Thus by Proposition 2.3 $H$ can be expressed as $H = (I - E - E)$. Solving for $E$ get $E = \frac{1}{2} \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}$. Thus $H = (I - E + e^{i\pi}E) = (I - E - E) = (I - 2E)$ expresses $H$ as a product of (minimal) basic unitary matrices where $E$ is as above. $H$ itself is a minimal basic matrix.

Then \{\$E, F\$\} with $F = I - E$ is a complete orthogonal set of idempotents. Now $HE = -E$ and $HF = F$ so $H$ ‘reverses’ $E$ and leaves $F$ alone.

Roots of $H$ are immediate: For example a square root of $H$ is $(I - E + e^{i\frac{\pi}{2}}E) = (I - E - iE)$. All the square roots of $H$ may be obtained by methods of Section 3.2 and these are $(I - E + e^{i\frac{\pi}{2}}E)(I - F + e^{i\pi}F)$ where $r \in \{1, 3\}, s \in \{0, 1\}$. The action of $H$ on the idempotents is given by $HF = F, HE = e^{i\pi}E; F$ is left alone by $H$ and $E$ is changed to $e^{i\pi}E$.

- **Pauli gates**: The Pauli gates are symmetric unitary matrices and hence satisfy $X^2 = I$ and have eigenvalues \{1, -1\}. Thus by Proposition 2.3, they have the form $(I - E + e^{i\pi}E) = I - 2E$ for an idempotent $E$. It is now a matter of establishing $E$ in these cases.

  - **Pauli X-gate** is $X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$ and has eigenvalues \{1, -1\}. Thus $I - 2E = X$ for a symmetric idempotent $E$. Solving for $E$ gives giving that $E = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & -\frac{1}{2} \end{pmatrix}$. Hence $X = (I - E + e^{i\pi}E)$.

  In this form a square root of $X$ is easy to find and one such is $(I - E + e^{i\frac{\pi}{2}}E)$. By section 3.2
all square roots of $X$ are $\{(I-E+e^{i\frac{r}{s}}E)(I-F+e^{i\frac{r}{s}}F)\}$ where $r \in \{1,3\}$, $s \in \{0,1\}$ and $F = I - E = \begin{pmatrix} 1 & -1 \\ 1 & 1 \end{pmatrix}$.

**Pauli Y-gate** is $Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$. Then $Y = I - E + e^{i\pi}E$ where $E = \begin{pmatrix} 1/2 & 1/2 \\ 1/2 & 1/2 \end{pmatrix}$.

Roots of Y-gate are easy to obtain: for example $(I-E+e^{i\pi}E)$ is a fourth root and all fourth roots may be directly found by structures of Section 3.2.

**Pauli Z-gate** is $Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$. Then $Z = (I-E+e^{i\pi}E) = I-2E$ where $E = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$.

The Pauli gates are themselves (minimal) basic matrices.

- **Phase shift gate**: The Phase shift unitary matrix is $Ph = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\theta} \end{pmatrix}$ and has eigenvalues $\{1, e^{i\theta}\}$.

  Applying the general result gives $Ph = (I-F+e^{i\theta}F)$ where $F = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$.

  This expresses $Ph$ as a minimal basic unitary matrix. Roots of $Ph$ are easily obtained by the method of section 3.2. For example a third root of $Ph$ is $(I-F+e^{i\theta}F)$ but there are others, see Section 3.2.

- **NOT and Square root of NOT**: The NOT gate is $NOT = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$. This has eigenvalues $\{1, -1\}$. Then $NOT = (I-E+e^{i\pi}E)$, where $E = \frac{1}{2} \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}$.

  Then (one value) of $\sqrt{NOT} = I - E + e^{i\pi}E = \frac{1}{2} \begin{pmatrix} 1 + i & 1 - i \\ 1 - i & 1 + i \end{pmatrix}$ for example it is possible to define $\sqrt{\sqrt{NOT}} = I - E + e^{i\pi}E$.

- **Swap, Square root of Swop**: The Swap gate has matrix $Sw = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$. This has eigenvalues $\lambda = 1$, (three times) and $\lambda = -1$. Then with $E = \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$ it is seen that $Sw = I - E + e^{i\pi}E$.

  Then a square root is $\sqrt{Sw} = I - E + iE = I - E + e^{i\pi}E$, where $E$ is the idempotent obtained for $Sw$ and this is $\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$ as usually given.

  $(I - E + e^{i\pi}E)$ could be considered as ‘a square root of a square root of Swap’.

- **CNOT matrix**: CNOT matrix is $CNOT = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$. This has eigenvalues 1, three times, and eigenvalue $-1$. Let $E = \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & -1 & -1 \\ 0 & 0 & -1 & -1 \end{pmatrix}$. Then $CNOT = (I-E+e^{i\pi}E)$. $E$ has rank 1 and $F = I - E$, corresponding to eigenvalue 1, has rank 3. All the square roots of CNOT are $\{(I-E+e^{i\pi}E),(I-E+e^{i\pi}E),(I-E+e^{i\pi}E)(I-F+e^{i\pi}F),(I-E+e^{i\pi}E)(I-F+e^{i\pi}F)\}$.

- **Bell non-symmetric**: The Bell non-symmetric matrix is $B = \begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}$. This has eigenvalues $\{i, -i\}$. Then it is easy to show that $B = (I-E-iE)(I-F+iF)$ where $E = \frac{1}{2} \begin{pmatrix} 1 & i \\ -i & 1 \end{pmatrix}$. $F = \frac{1}{2} \begin{pmatrix} 1 & -i \\ i & 1 \end{pmatrix} = I - E$. Thus $B = (I-E-iE)(I-F+iF) = (I-E+e^{i\pi}E)(I-F+e^{i\pi}F)$. Note that $\{E, F\}$ is complete.
5 Build matrices

Normal matrices are built from basic matrices of the form \((I - E + \alpha E)\); unitary matrices or quantum logic gates are built from basic matrices of the form \((I - E + e^{i\theta} E)\). A particular property may be required; for example a quantum logic gate which is an \(n^{th}\) root and/or a quantum logic gate with particular eigenvalues and multiplicities may be required. Requirements may be met in a constructive manner.

It is shown in Proposition 4.6 that a symmetric (unitary matrix)/(quantum logic gate) is of the form \((I - 2E)\) for a symmetric idempotent \(E\). This makes building symmetric unitary matrices/logic gates particularly easy.

For example the matrix \(U = \frac{1}{2} \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{pmatrix} \) is the unitary matrix related to the Hadamard matrix \(2U\). \(U\) is symmetric and so \(U = (I - 2E)\) where \(E\) is easily shown to be \(E = \frac{1}{4} \begin{pmatrix} 1 & -1 & -1 & 1 \\ -1 & 1 & 1 & -1 \\ 1 & -1 & 1 & 1 \\ -1 & 1 & -1 & 1 \end{pmatrix} \). It is clear that \(\text{rank}(E) = 2\) and thus \(\text{rank}(I - 2E) = 2\). A square root of \(U\) is \((I + E + e^{i\pi} E)\) but there are three other square roots which are easily worked out using the methods developed in Section 3.2. \(U\) is not separable and roots of \(U\) are not separable.

A unitary matrix \(U\) which also is a root of \(I\) may be built by the methods of Section 3.2 in this case \(U^n = I\) and \(U^* = U^{n-1}\). Suppose for example quantum logic gates (unitary matrices) in \(C_{4 \times 4}\) are required which are \(4^{th}\) roots of \(I_4 = I\). Clearly \(U = (I - E + e^{i \frac{2\pi}{4}} E)\) where \(E\) is any idempotent and \(r \in \{1,2,3\}\), is a basic unitary quantum gate which is a \(4^{th}\) root of \(I\). Now \(v_1 = \frac{1}{2}(1,1,-1,1)^T, v_2 = \frac{1}{2}(1,-1,1,1)^T, v_3 = \frac{1}{2}(-1,1,1,1)^T\) is an orthonormal basis for \(E_1\) and these give an orthogonal complete set of idempotents \(E_1 = v_1 v_1^*, E_2 = v_2 v_2^*, E_3 = v_3 v_3^*, E_4 = v_4 v_4^*\). From these, gates which are \(4^{th}\) roots of \(I\) may be formed: \(\{ (I - E_1 + e^{i \frac{2\pi}{4}} E_1)(I - E_2 + e^{i \frac{2\pi}{4}} E_2)(I - E_3 + e^{i \frac{2\pi}{4}} E_3)(I - E_4 + e^{i \frac{2\pi}{4}} E_4) \} \) where \(r_j \in \{0,1,2,3\}\). (If all \(r_j = 0\) then the identity matrix is obtained.)

Normal including unitary and Hermitian matrices may naturally be built from a set of unit orthogonal vectors. Let \(\{u_1, u_2, \ldots, u_k\}\) be a set of orthogonal unit column vectors in \(C_n\) and set \(E_i = u_i u_i^*\). Build \(B = (I - E_1 + \alpha_1 E_1)(I - E_2 + \alpha_2 E_2) \ldots (I - E_k + \alpha_k E_k)\) which is then normal. The basic matrices with the same scalars may be amalgamated by Lemma 2.4 to get a unique expression \(B = (I - F_1 + \beta_1 F_1)(I - F_2 + \beta_2 F_2) \ldots (I - F_s + \beta_s F_s)\) where the \(\beta_i\) are distinct. Each \(F_j\) is a sum of \(\alpha_i u_i u_i^*\) of a subset of \(\{u_1, u_2, \ldots, u_k\}\). Note that \(S = \{u_1, u_2, \ldots, u_k\}\) need not be a full basis for \(C_n\); in this case \(\{E_1, E_2, \ldots, E_k\}\) is not a complete set of orthogonal idempotents. By letting \(E = I - \sum_{i=1}^k E_i\) a complete set \(\{S \cup E\}\) of orthogonal idempotents is obtained. To build a unitary matrix take \(\alpha_i = e^{i \theta_i}\) and to build a Hermitian matrix take the \(\alpha_i\) to be real. If no \(\alpha_i = 0\) then no \(\beta_j = 0\) and \(B\) is invertible with inverse \(B^{-1} = (I - F_1 - \beta_1^{-1} F_1)(I - F_2 - \beta_2^{-1} F_2) \ldots (I - F_s - \beta_s^{-1} F_s)\). If some \(\beta_j = 0\), say \(b_k = 0\), then \(B = (I - F_1 - \beta_1^{-1} F_1)(I - F_2 - \beta_2^{-1} F_2) \ldots (I - F_{s-1} + \beta_{s-1} F_{s-1}) (I - F_s)\). In this case the pseudo inverse of \(B\) is \((I - F_1 - \beta_1^{-1} F_1)(I - F_2 - \beta_2^{-1} F_2) \ldots (I - F_{s-1} - \beta_{s-1} F_{s-1}) (I - F_s)\). Sets of orthogonal unit column \(n \times 1\) vectors may be obtained from \(n^{th}\) roots of unity from which unitary matrices and logic gates for various requirements may be built; details are omitted.

Comparisons may be made with the famous 1D unique building blocks for paraunitary matrices due to Belvitch and Vaidyanathan, see [5] pages 302-322.

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