QUASIMONTE CARLO METHODS AND PSEUDO-RANDOM NUMBERS

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Nothing in Nature is random. . . . A thing appears random only through the incompleteness of our knowledge.

Spinoza, Ethics I

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1. Introduction. The subject matter of this talk is at the crossroads of two areas which will turn out to have more than only an etymological kinship, namely numerical analysis and number theory. Like so many mixed breeds, it has its fascinations and attractions, but also its inherent dilemmas. A multitude of concepts and devices dear to numerical analysts and computer users are, in open or disguised form, of an arithmetic nature, and problems arising in the computational workshop, especially those requiring effective methods, are now treated quite frequently with the powerful tools of the number theorist. This provides for a vivid interplay and is a source of enrichment for both disciplines. Of course, the occasion only permits us to look at a certain segment in the broad spectrum of activities. The leitmotif in our discussion will be the simulation of procedures containing an element of randomness by judiciously chosen deterministic schemes, with number theory playing a
prominent part in the construction of such schemes.

Our exposition can be roughly divided into two parts, which will not preclude, however, some strong interrelations between these. §§2–5 are devoted to deterministic versions of Monte Carlo techniques that have come to be known under the collective term “quasi-Monte Carlo methods”. The widest range of applications, and indeed the historical origin of these methods, is found in numerical integration, but related matters such as interpolation problems and the numerical solution of integral equations can also be dealt with successfully. It does not seem to be too well known among the craftsmen of the trade that quasi-Monte Carlo methods possess two big assets not shared by standard Monte Carlo techniques, namely effectiveness and fast convergence. It is therefore hoped that this talk will help in the dissemination and eventual adoption of these more efficient methods.

As we already indicated, the term “quasi-Monte Carlo” encompasses a variety of techniques. In their basic form, they all emerged sometime in the 1950s, and another common bond between them is their heavy reliance on number-theoretic concepts. These methods fall into three main categories which are characterized by the titles of §§3–5. For the sake of completeness, we will present the (at least from the standpoint of the specialist) classical foundations of these various techniques, but emphasize the refinements and extensions in scope achieved in recent years. This vigorous progress has been sustained by significant contributions from both the Russian school and the Western branch, with the Chinese doing more than just checking the balance.

In the second part, comprising §§6–11, we are concerned with the vital matter of pseudo-random number generation. The limelight will be on the most popular generators, namely Lehmer’s linear congruential pseudo-random numbers. These appeared to fall into disfavor among theoreticians about ten years ago upon the disclosure of a certain undesirable lattice (or “crystalline”) structure, but there is every indication that the practitioners did not care much about these squabbles and continued defiantly with their time-honored routines. Recent results of the author amount to a rehabilitation of these generators and justify a posteriori the woman (or man) at the computer in her (or his) preservative attitude. To sum up a complex matter succinctly, this research has shown that a proper choice of parameters in the generation procedure leads to a sequence of pseudo-random numbers enjoying excellent properties of statistical (almost-) independence among a given number of successive terms. The italicized part of the preceding sentence cannot be emphasized too much. Although not totally ignored, the matter of adequately selecting these parameters has all too frequently been left to chance or was based on insufficient evidence. If, as so often, generators with the said statistical independence properties are required, we can now provide for the first time reliable criteria for the choice of parameters which are buttressed by effective theoretical results.

Lest a contrary impression be created, it should be pointed out that the existence of linear congruential pseudo-random numbers with desirable statistical independence properties does not negate the results about their unfavorable lattice structure. The fact is simply that the lattice structure does not detract from the usefulness of these generators as long as we use the
pseudo-random numbers for purposes in which only the statistical independence of successive terms or a good distribution behavior are relevant. Whoever finds that there is still an irreconcilable dilemma here, may want to ponder the following morphological principle: if one attempts to distribute points in a well-planned and equitable manner over a given domain, one will, intentionally or inadvertently, provide the resulting collection of points with an intrinsic structure. Consider a very simple example in which we are challenged to distribute denumerably many points on the real axis in what we deem the fairest way. We will most likely end up with an equally spaced arrangement and have thereby generated a point set with the structure of a one-dimensional lattice. Thus, a lattice structure may even be a virtue when good distribution properties are desired. More will be said about this in §§10 and 11.

The advance of our knowledge in this area of linear congruential pseudo-random numbers has gone hand in hand with progress in pure number theory, namely the establishment of nontrivial estimates for exponential sums with linear recurring arguments. We have found it convenient to relegate these number-theoretic matters to a separate section (see §8). Apart from the topics already mentioned, we shall discuss related work on Lehmer’s pseudo-random numbers and similar generators as well as review their elementary properties.

As to the list of references, we have attempted to be fairly complete concerning the literature on quasi-Monte Carlo methods since no comprehensive survey of this area was available up to now. With respect to pseudo-random numbers, the prolific output in this discipline has been assessed periodically in review articles and bibliographies (cf. [84], [100], [135], [207], [208], [322a]) and there are monographs covering the subject (cf. [80], [142], [154], [182], [193], [213]), so that we have only listed those works having a direct bearing on our discussion.

Before we embark on our exposition proper, it is necessary to gain an understanding of the fundamental principle involved in the Monte Carlo method which is the progenitor of the great bulk of the theory to be expounded here. The Monte Carlo method (or “method of statistical trials”) may be described in simple terms as a numerical method based on random sampling. We give an illustration below. The history of the method has been charted often enough. Its birth is assumed to have taken place in 1949 with the publication of [199], although it was definitely known earlier to a clandestine group working on U. S. Defense projects. Statisticians have intuitively used its principle long before that (cf. [334]), but it was only the computer age that could turn it into a systematic and viable technique. A number of excellent monographs explore the whole range of the method (cf. [306], p. 9) points out correctly, the method is not of much help in trying to win at roulette.

Mainly J. von Neumann, N. Metropolis, S. M. Ulam, H. Kahn, and their collaborators at the Los Alamos Scientific Laboratory.

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2 Most applications in numerical analysis are of this type.
3 A Rorschach test could be set up interpreting deviations from this arrangement in psychological terms.
4 As Sobol' [306, p. 9] points out correctly, the method is not of much help in trying to win at roulette.
5 Mainly J. von Neumann, N. Metropolis, S. M. Ulam, H. Kahn, and their collaborators at the Los Alamos Scientific Laboratory.
We also recommend the survey article \[100\], the recent bibliography \[84\], and the elementary introduction in \[306\]. The main reason for the popularity of the Monte Carlo method is its applicability to a never-ending variety of problems in numerical analysis, statistics, applied mathematics, particle physics, engineering, systems analysis, and so on. We refer to \[28\], \[29\], \[70\], \[83\], \[106\], \[151\], \[200\], \[323\], \[340\] for accounts of some of these applications. For the general area of simulation, see \[17\], \[27\], \[69\], \[79\], \[190\], \[201\], \[209\].

To present a simple example of a Monte Carlo calculation, we consider the problem of computing the area of a region $E$ of complicated shape contained in the unit square $[0, 1] \times [0, 1]$. The idea is now to select at random $N$ points from the unit square by performing $N$ independent trials. In practice, $N$ should be fairly large, say $N \approx 10^4$. If $x_1, \ldots, x_N$ are the points resulting from the sampling process, then the Monte Carlo approximation is

\[
\text{area of } E \approx \frac{1}{N} \sum_{n=1}^{N} c_E(x_n),
\]

where $c_E$ is the characteristic function of $E$. In other words, the fraction of the sample points falling into $E$ is taken as an approximate value for the area of $E$. This procedure can be generalized immediately if one recognizes the left-hand side of (1.1) as the (Lebesgue) integral of $c_E$ over the unit square (provided that $E$ is Lebesgue-measurable, of course). Thus, for a given dimension $s > 1$, let $I^s = [0, 1]^s$ be the $s$-dimensional unit cube and let $f = f(t)$ be a bounded \footnote{This condition is adopted to avoid technicalities and can be relaxed.} Lebesgue-integrable function on $I^s$. Then the Monte Carlo approximation for the Lebesgue integral of $f$ over $I^s$ is

\[
\int_{I^s} f(t) \, dt \approx \frac{1}{N} \sum_{n=1}^{N} f(x_n),
\]

where $x_1, \ldots, x_N$ are random points from $I^s$ obtained by $N$ independent trials. Therefore, the basic principle of integration by the Monte Carlo method is to replace a continuous average by a discrete average over randomly selected points. In the same vein, if $E$ is a Lebesgue-measurable subset of $I^s$, then we may say on the basis of (1.2) that

\[
\int_{E} f(t) \, dt = \int_{I^s} f(t)c_E(t) \, dt \approx \frac{1}{N} \sum_{n=1}^{N} f(x_n)c_E(x_n),
\]

and so the Monte Carlo approximation is taken to be

\[
\int_{E} f(t) \, dt \approx \frac{1}{N} \sum_{n=1}^{N} f(x_n),
\]

where $x_1, \ldots, x_N$ are as in (1.2).

The strong law of large numbers guarantees that the numerical integration procedure (1.2) converges almost surely. Moreover, it follows from the central limit theorem that the expected integration error is $O(N^{-1/2})$. The remarkable feature here is that this order of magnitude does not depend on the
dimension. This explains the interest in the Monte Carlo method for large
dimensions, where classical techniques perform poorly. It should be
mentioned, however, that the implied constant in this estimate depends on a
certain variance factor through which the dimension may enter. The idea of
reducing this variance by suitable transformations plays an important role in
the Monte Carlo method (cf. [66], [100]).

For the practical implementation of the Monte Carlo method, the funda­
mental question is, of course, how to produce a random sample. There is no
ready-made answer since no satisfactory definition of randomness exists (see
§6 for an elaboration on this point). Some people use tables of “random”
numbers such as [251] or physical devices for generating random numbers
such as white noise. But there is now an ever expanding school of thought
which has come to realize that instead of trying to cope with the impalpable
concept of randomness, one should select points according to a deterministic
scheme that is well suited for the problem at hand. This is the underlying idea
of a quasi-Monte Carlo method. For instance, in the area of numerical
integration it turns out to be quite irrelevant whether the sample points or
“nodes” are truly random; of primary importance is really the even distribu­
tion of the points over $I^s$ (compare with §2). Thus, rather than worrying
about random selection procedures, one should be concerned with finding
sets of nodes having an optimally fair distribution (see §3).

If we want to emphasize the distinction between quasi-Monte Carlo meth­
ods and the standard Monte Carlo method, we will employ the term
“statistical Monte Carlo method” for the latter. There is another differen­
tiation in language that will occur in the course of our discussion and that is
sometimes regarded as artificial, namely that between quasi-random points
(or numbers) and pseudo-random numbers. Although no clear-cut line can be
drawn, there is a subtle distinction here, in the sense that the use of
quasi-random points is customarily restricted to numerical integration or very
closely related applications and that, consequently, they only have to show an
acceptable distribution behavior, whereas pseudo-random numbers are
supposed to serve a multitude of purposes and should therefore perform well
under a battery of statistical tests.

PART I. QUASI-MONTE CARLO METHODS

2. Quasi-Monte Carlo integration. We noted already that numerical
integration by a quasi-Monte Carlo method depends on the judicious choice
of nodes from the integration domain or a superset thereof. This poses the
question as to the precise criteria according to which the nodes should be
selected. To find out, let us first consider the case where the integration
domain is $I^s = [0, 1]^s$. Here we use the Monte Carlo approximation

$$\int_{I^s} f(t) \, dt \approx \frac{1}{N} \sum_{n=1}^{N} f(x_n). \quad (2.1)$$

For the sake of this discussion, we adopt a simpler model by replacing the
large set of nodes $x_1, \ldots, x_N$ by an infinite sequence $x_1, x_2, \ldots$ of points in
$I^s$. Then as we increase $N$ in (2.1), we obviously want the integration error to
become negligible. Thus we require that

$$\lim_{N \to \infty} \frac{1}{N} \sum_{n=1}^{N} f(x_n) = \int_{I^s} f(t) \, dt.$$  \hspace{1cm} (2.2)

This limit relation should hold for a reasonable class of integrands, say for all continuous functions $f$ on $I^s$. The resulting condition on the sequence $x_1, x_2, \ldots$ is precisely one of the well-known criteria for this sequence to be uniformly distributed in $I^s$. This concept is usually defined as follows. The sequence $x_1, x_2, \ldots$ of points in $I^s$ is called uniformly distributed in $I^s$ if

$$\lim_{N \to \infty} \frac{1}{N} \sum_{n=1}^{N} c_f(x_n) = |J|$$

holds for all subintervals $J$ of $I^s$, where $|J|$ denotes the $s$-dimensional Lebesgue measure (= volume) of $J$. Intuitively, this means that the points $x_1, x_2, \ldots$ are spread out over the unit cube $I^s$ according to the principle of proportional representation. A detailed treatment of uniformly distributed sequences can be found in the book of Kuipers and Niederreiter [174].

In case the sequence $x_1, x_2, \ldots$ is uniformly distributed in $I^s$, the relation (2.2) actually holds for all Riemann-integrable functions $f$ on $I^s$ (cf. [174, pp. 3, 52]). On the other hand, (2.2) need not hold for arbitrary Lebesgue-integrable functions; e.g., it fails if $f$ is the characteristic function of the set $\{x_1, x_2, \ldots \}$. In fact, (2.2) characterizes Riemann-integrability in the following sense: if $f$ is a real-valued function on $I^s$ such that the limit in (2.2) exists for all uniformly distributed sequences in $I^s$, then $f$ must be Riemann-integrable on $I^s$ (de Bruijn and Post [54], Binder [22]). Therefore, numerical integration by a quasi-Monte Carlo technique should only be employed for a Riemann-integrable integrand since only in this case can we guarantee convergence. On the theoretical level, this is a significant difference as compared to a statistical Monte Carlo method, for which the strong law of large numbers affirms the almost sure convergence in (2.2) for any bounded Lebesgue-integrable $f$. But for practical purposes, the restriction to Riemann-integrable functions in quasi-Monte Carlo methods is not of a serious nature.

We turn now to the general case of an integration domain $E \subseteq I^s$. As we have seen in (1.3), the Monte Carlo approximation attains the form

$$\int_{E} f(t) \, dt \approx \frac{1}{N} \sum_{n=1}^{N} f(x_n).$$

If we use again an infinite sequence as a model, then by what we have already learned, the sequence $x_1, x_2, \ldots$ should be uniformly distributed in $I^s$. Further inspection shows that the integration domain $E$ cannot be quite arbitrary if we want convergence of the method. For even if we take a simple integrand such as the constant function $f \equiv 1$ (i.e., if we are asked to calculate the volume of $E$), we arrive at the convergence condition

$$\lim_{N \to \infty} \frac{1}{N} \sum_{n=1}^{N} c_E(x_n) = \int_{I^s} c_E(t) \, dt,$$

which need only hold if $c_E$ is Riemann-integrable, or, equivalently, if $E$ itself
is Jordan-measurable (= has an elementary volume). In toto, we get convergence in a quasi-Monte Carlo integration procedure if the integrand \( f \) is Riemann-integrable, the integration domain \( E \subseteq I^s \) is Jordan-measurable, and the sequence \( x_1, x_2, \ldots \) of nodes is uniformly distributed in \( I^s \).

Returning from the question of convergence to the original setup, namely that of a finite collection of nodes, we realize that such a discrete set can never constitute a completely fair distribution over \( I^s \) since there will always be subintervals \( J \) of \( I^s \) (possibly of very small volume) which do not contain any one of the given points. Thus, the uniform distribution property is an idealization, and in actual practice we have to settle for an approximation. The model of an infinite sequence has provided the clue to the proper criterion for selecting nodes, viz. the even distribution of the nodes over \( I^s \). Therefore, we shall introduce a quantity which measures the uniformity of distribution of a given set of nodes.

We consider first the one-dimensional case. Let \( x_1, \ldots, x_N \) be \( N \) numbers in \( I = [0, 1] \). If \( E \) is a subset of \( I \), then

\[
A(E; N) = \sum_{n=1}^{N} c_E(x_n)
\]

counts the number of \( n, 1 < n < N \), with \( x_n \in E \).

2.1. Definition. The discrepancy \( D_N \) of the \( N \) numbers \( x_1, \ldots, x_N \) in \( I \) is defined by

\[
D_N = \sup_J \left| \frac{A(J; N)}{N} - \frac{1}{|J|} \right|,
\]

where \( J \) runs through all subintervals of \( I \) and \( |J| \) denotes the length of \( J \).

It is immaterial whether one considers closed, open, half-open, or arbitrary intervals \( J \) since any one category leads to the same value of the supremum in (2.3), as can be seen from [174, p. 99]. A useful variant of the above definition is the following.

2.2. Definition. The discrepancy \( D_N^* \) of the \( N \) numbers \( x_1, \ldots, x_N \) in \( I \) is defined by

\[
D_N^* = \sup_{0 < t < 1} \left| \frac{A([0, t); N)}{N} - t \right|.
\]

Finite sequences of nodes with small discrepancy \( D_N \) resp. \( D_N^* \) provide a valid approximation to uniform distribution, in the sense that for an infinite sequence both \( \lim_{N \to \infty} D_N = 0 \) and \( \lim_{N \to \infty} D_N^* = 0 \) are equivalent to the sequence being uniformly distributed in \( I \) (cf. [174, Chapter 2, §1]). Here \( D_N \) resp. \( D_N^* \) stands for the discrepancy of the first \( N \) terms of the sequence.

In nonparametric statistics the method of measuring the deviation between the empirical distribution of \( x_1, \ldots, x_N \) and the uniform distribution on \( I \) by the quantity \( D_N^* \) is also known as the two-sided Kolmogorov test. The discrepancy \( D_N^* \) may also be thought of as the supremum norm of the function

\[
R_N(t) = N^{-1} A([0, t); N) - t, \quad 0 < t < 1.
\]
By taking various other norms of this function, one arrives at further concepts of discrepancy, the most commonly used among these being the $L^2$ discrepancy $T_N$ given by

$$T_N = \left( \int_0^1 R_N(t)^2 \, dt \right)^{1/2}.$$ 

See [220] for further information about the $L^2$ discrepancy. One may also consider discrepancies with respect to distribution functions different from the uniform distribution ([125], [219]) and with respect to other summation methods ([121], [222], [223]).

The significance of the discrepancy stems from the fact that it occurs in the effective error estimates for quasi-Monte Carlo integration. The precise form of these estimates depends on the regularity of the integrand.

2.3. **Theorem (Koksma [156]).** If $f$ is a function of bounded variation $V(f)$ on $I$ and $x_1, \ldots, x_N$ are numbers in $I$ with discrepancy $D^*_N$, then

$$\left| \frac{1}{N} \sum_{n=1}^N f(x_n) - \int_0^1 f(t) \, dt \right| \leq V(f)D^*_N. \quad (2.4)$$

This error bound is particularly appealing since the influences on the integration error are clearly separated: the regularity of the integrand is reflected in the factor $V(f)$ and the uniformity of distribution of the nodes is controlled by the discrepancy $D^*_N$.

If $f$ is of bounded variation and continuous on $I$, then Koksma’s inequality can be proved very quickly using integration by parts. For

$$\int_0^1 R_N(t) \, df(t) = \frac{1}{N} \sum_{n=1}^N \int_0^1 c_{0,1}(x_n) \, df(t) - \int_0^1 t \, df(t)$$

$$= \frac{1}{N} \sum_{n=1}^N (f(1) - f(x_n)) - f(1) + \int_0^1 f(t) \, dt$$

$$= -\frac{1}{N} \sum_{n=1}^N f(x_n) + \int_0^1 f(t) \, dt,$$

and so

$$\left| \frac{1}{N} \sum_{n=1}^N f(x_n) - \int_0^1 f(t) \, dt \right| = \left| \int_0^1 R_N(t) \, df(t) \right|$$

$$< V(f) \sup_{0 \leq t \leq 1} |R_N(t)| = V(f)D^*_N.$$ 

The general case is shown by a slight variation of this argument (see [174, p. 143]).

2.4. **Theorem (Niederreiter [217]).** If $f$ is a continuous function on $I$ with modulus of continuity $M$ and $x_1, \ldots, x_N$ are numbers in $I$ with discrepancy $D^*_N$, then

$$\left| \frac{1}{N} \sum_{n=1}^N f(x_n) - \int_0^1 f(t) \, dt \right|$$

$$< V(f) \sup_{0 \leq t \leq 1} |R_N(t)| = V(f)D^*_N.$$ 

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QUASI-MONTE CARLO METHODS

if $f(\cdot) - f(0) \leq M(D_N^*)$.

A general error estimate valid for any Riemann-integrable integrand was established by Hlawka [120]. Analogues of the above inequalities can be shown for other summation methods ([222], [223, pp. 148-149], [303, Chapter 2]) and for other distribution functions ([110], [125], [307]). Estimates in terms of the $L^2$ discrepancy are also available ([292], [295], [366]) and are obtained by applying the Cauchy-Schwarz inequality at a certain stage in the argument. For instance, if $f$ has a continuous derivative on $I$ and $x_0, \ldots, x_N$ are numbers in $I$ with $L^2$ discrepancy $T_N$, then

$$\left| \frac{1}{N} \sum_{n=1}^{N} f(x_n) - \int_0^1 f(t) \, dt \right| \leq \left( \int_0^1 (f'(t))^2 \, dt \right)^{1/2} T_N.$$ 

Error bounds of a different nature result from measuring the regularity of a periodic integrand by the size of its Fourier coefficients ([140], [214]).

We turn now to the multidimensional case in which quasi-Monte Carlo methods are usually applied. Let $I^s = [0, 1]^s$ be the $s$-dimensional unit cube. For $N$ given points $x_1, \ldots, x_N$ in $I^s$ and a subset $E$ of $I^s$, we introduce the counting function

$$A(E; N) = \sum_{n=1}^{N} c_E(x_n).$$

We use $|E|$ to denote the $s$-dimensional Lebesgue measure of $E$. To unify various definitions of discrepancy that we shall need, we start from a general concept.

2.5. DEFINITION. Let $\mathfrak{M}$ be a nonempty family of Lebesgue-measurable subsets of $I^s$. Then the discrepancy $D_N(\mathfrak{M})$ of the $N$ points $x_1, \ldots, x_N$ in $I^s$ is defined by

$$D_N(\mathfrak{M}) = \sup_{E \in \mathfrak{M}} \left| \frac{A(E; N)}{N} - |E| \right|.$$ 

2.6. DEFINITION. The discrepancy $D_N$ of the $N$ points $x_1, \ldots, x_N$ in $I^s$ is defined by $D_N = D_N(\mathfrak{J})$, where $\mathfrak{J}$ is the family of all subintervals of $I^s$, and the discrepancy $D_N^*$ is defined by $D_N^* = D_N(\mathfrak{J}^*)$, where $\mathfrak{J}^*$ is the family of subintervals of $I^s$ of the form $[0, t_1] \times \cdots \times [0, t_s]$.

For an infinite sequence of points in $I^s$, we have as in the one-dimensional case that both $\lim_{N \to \infty} D_N = 0$ and $\lim_{N \to \infty} D_N^* = 0$ are equivalent to the sequence being uniformly distributed in $I^s$ (cf. [174, Chapter 2, §1]).

In order to generalize Koksma's inequality, we have to set up an appropriate concept of total variation for functions of several variables. For a function $f$ on $I^s$ and an interval $J = [a_1^{(1)}, a_2^{(1)}] \times \cdots \times [a_1^{(s)}, a_2^{(s)}] \subseteq I^s$, we put

$$\text{More generally, the integration error can be estimated in terms of an appropriately defined } L^p \text{ discrepancy [303, Chapter 2].}$$

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\[ \Delta(f; J) = \sum_{\varepsilon_1 = 1}^{2} \cdots \sum_{\varepsilon_s = 1}^{2} (-1)^{\varepsilon_1 + \cdots + \varepsilon_s} f(a^{(1)}_1, \ldots, a^{(s)}_s). \]

To define a partition \( \mathcal{P} \) of \( I^s \), we start from \( s \) finite sequences of the form \( 0 = \eta^{(j)}_0 < \eta^{(j)}_1 < \cdots < \eta^{(j)}_{m_j} = 1 \) (\( j = 1, 2, \ldots, s \)). The partition then consists of all the intervals \([\eta^{(j)}_i, \eta^{(j)}_{i+1}] \times \cdots \times [\eta^{(s)}_i, \eta^{(s)}_{i+1}]\) with \( 0 < i_j < m_j \) for \( j = 1, 2, \ldots, s \).

2.7. Definition. For a function \( f \) on \( I^t \), we set

\[ V^{(s)}(f) = \sup_{\mathcal{P}} \sum_{J \in \mathcal{P}} |\Delta(f; J)|, \]

where the supremum is extended over all partitions \( \mathcal{P} \) of \( I^s \). If \( V^{(s)}(f) \) is finite, then \( f \) is said to be of bounded variation on \( I^s \) in the sense of Vitali.

For functions \( f = f(t_1, \ldots, t_s) \) that are sufficiently regular, \( V^{(s)}(f) \) can be represented by an integral; namely,

\[ V^{(s)}(f) = \int_0^1 \cdots \int_0^1 \left| \frac{\partial f}{\partial t_1} \cdots \frac{\partial f}{\partial t_s} \right| dt_1 \cdots dt_s \quad (2.5) \]

whenever the indicated partial derivative is continuous on \( I^s \). If \( f \) actually depends on less than \( s \) variables, then we always have \( \Delta(f; J) = 0 \), and so \( V^{(s)}(f) = 0 \). Since such a function \( f \) may still be highly irregular, we have to consider a more suitable notion of variation which is obtained by taking into account the behavior of \( f \) on the various faces of \( I^s \).

2.8. Definition. Let \( f \) be a function on \( I^s \). For \( 1 < k < s \) and \( 1 < i_1 < i_2 < \cdots < i_k < s \), we denote by \( V^{(k)}(f; i_1, \ldots, i_k) \) the \( k \)-dimensional variation in the sense of Vitali of the restriction of \( f \) to \( I_{i_1} \cdots i_k = \{ (t_1, \ldots, t_s) \in I^s : t_j = 1 \text{ for } j \neq i_1, \ldots, i_k \} \). If all variations \( V^{(k)}(f; i_1, \ldots, i_k) \) are finite, then \( f \) is said to be of bounded variation on \( I^s \) in the sense of Hardy and Krause.

Since \( V^{(s)}(f) = V^{(s)}(f; 1, 2, \ldots, s) \), a function of bounded variation in the sense of Hardy and Krause is automatically of bounded variation in the sense of Vitali.

2.9. Theorem (Hlawka [111]). If \( f \) is a function of bounded variation on \( I^s \) in the sense of Hardy and Krause and \( x_1, \ldots, x_N \) are points in \( I^s \), then

\[ \left| \frac{1}{N} \sum_{n=1}^{N} f(x_n) - \int_{I^s} f(t) \, dt \right| < \sum_{k=1}^{s} \sum_{1 < i_1 < i_2 < \cdots < i_k < s} V^{(k)}(f; i_1, \ldots, i_k) D_N^* (i_1, \ldots, i_k), \quad (2.6) \]

where \( D_N^*(i_1, \ldots, i_k) \) is the discrepancy in \( I_{i_1} \cdots i_k \) of the points obtained by orthogonal projection of \( x_1, \ldots, x_N \) onto \( I_{i_1} \cdots i_k \).

The above estimate is often called the Koksma-Hlawka inequality. A proof of this result based on multidimensional integration by parts was given.

Information about these concepts of variation can be found in Hobson [127] and the literature in [174, p. 158].

For sufficiently regular integrands, Hlawka [112] presents a simplified proof of the inequality, with the variations replaced by the corresponding integrals in (2.5).
by Zaremba [366] and is reproduced in [174, Chapter 2, §5].

One obtains a simplified form of (2.6) if one defines the variation \( V(f) \) of \( f \) on \( I^s \) in the sense of Hardy and Krause to be

\[
V(f) = \sum_{k=1}^{s} \sum_{1 < i_1 < i_2 < \cdots < i_k < s} V^{(k)}(f; i_1, \ldots, i_k).
\]

Obviously, \( V(f) \) is finite if and only if \( f \) is of bounded variation on \( I^s \) in the sense of Hardy and Krause. By using (2.5), \( V(f) \) can be written in terms of integrals for sufficiently regular \( f \). Since we always have

\[
D^*_N(i_1, \ldots, i_k) < D^*_N(1, 2, \ldots, s) = D^*_N,
\]

it follows from (2.6) that

\[
\left| \frac{1}{N} \sum_{n=1}^{N} f(x_n) - \int_{I^s} f(t) \, dt \right| < V(f)D^*_N.
\]

(2.7)

This inequality is completely analogous to (2.4).

The Koksma-Hlawka inequality can be generalized to more abstract settings ([214], [278]). Error estimates for other than equal-weight formulas are also known ([222], [223, pp. 148–149]). Inequalities for Riemann-integrable functions ([120]) and error bounds in terms of an appropriately defined \( L^2 \) discrepancy \( T^*_N \) ([292], [366], [101]) or an \( L^1 \) discrepancy ([292], [303, Chapter 8]), and more generally for an \( L^p \) discrepancy ([303, Chapter 8]), have been established as well. See [126] for the case of a nonuniform distribution function.

A measure for the evenness of distribution different from the discrepancy has also been considered in the literature. To introduce this concept, we need some preliminaries. By a dyadic interval, we mean a subinterval of \( I \) of the form \([j2^{-m}, (j + 1)2^{-m}]\) with integers \( m > 1 \) and \( 0 < j < m \), and with the stipulation that the interval be closed if its right endpoint is 1. In the \( s \)-dimensional case, a dyadic box is meant to be a cartesian product of any \( s \) dyadic intervals. Given a dyadic box \( B \) in \( I^s \), we split it into two parts as follows. For the moment, we move the origin of the coordinate system to the center of \( B \) and denote by \( \xi_1, \ldots, \xi_s \) the new coordinates. Then the union of those “quadrants” of \( B \) in which \( \text{sgn}(\xi_1 \cdots \xi_s) = (-1)^s \) is called the positive part \( B^+ \) of \( B \), whereas the remaining portion of \( B \) is the negative part \( B^- \). Let now \( x_1, \ldots, x_N \) be \( N \) given points in \( I^s \). Then we define the \( s \)-dimensional nonuniformity of these points to be

\[
\sup_B |A(B^+; N) - A(B^-; N)|,
\]

where the supremum is extended over all possible dyadic boxes \( B \) in \( I^s \). Furthermore, we project the given points orthogonally onto the various \( k \)-dimensional faces of \( I^s \) (\( 1 < k < s \)) and calculate the \( k \)-dimensional nonuniformity of the projected points in the respective face. Finally, the largest value among the \( s \)-dimensional nonuniformity and all these lower-dimensional nonuniformities is said to be the nonuniformity \( \varphi(\infty)(N) \) of the points \( x_1, \ldots, x_N \) (Sobol’ [291]). It is always a positive integer < \( N \). The motivation for this definition comes from the theory of Haar functions (compare with
For \( 1 < q < \infty \), an \( L^q \) analogue \( \varphi_q(N) \) has also been introduced (Sobol' \[289\]). If \( x_1, x_2, \ldots \) is a sequence of points in \( I^s \), then \( \varphi_q(N) \), \( 1 < q < \infty \), is the appropriate nonuniformity of the first \( N \) terms of the sequence. For any \( q, 1 < q < \infty \), we have the criterion that a sequence is uniformly distributed in \( I^s \) if and only if \( \lim_{N \to \infty} \varphi_q(N)/N = 0 \) \([289], [294]\). A comparison with the definition of discrepancy leads easily to the inequality

\[
\varphi_\infty(N) \leq 2^s ND_N \tag{2.8}
\]

for any \( N \) points in \( I^s \).

Estimates for integration errors can also be established in terms of nonuniformities. For instance, if \( f(t) = f(t_1, \ldots, t_s) \) is such that all its mixed partial derivatives

\[
\frac{\partial^k f}{\partial t_{i_1} \cdots \partial t_{i_k}}, \quad 1 \leq i_1 < i_2 < \cdots < i_k \leq s, \quad 1 \leq k \leq s,
\]

exist and are continuous on \( I^s \), then

\[
\left| \frac{1}{N} \sum_{n=1}^{N} f(x_n) - \int_{I^s} f(t) \, dt \right| \leq C(f) \frac{\varphi_\infty(N) \log^s N}{N}, \tag{2.9}
\]

where \( C(f) \) is a constant depending on \( f \) that reflects, as usual, the amount of oscillation of the integrand and \( \varphi_\infty(N) \) is the nonuniformity of the nodes \( x_1, \ldots, x_N \) (Sobol' \[291\], \[303, Chapter 4\]). Analogous estimates using the nonuniformities \( \varphi_q(N) \), \( 1 < q < \infty \), are also available \([289], [303, Chapter 4]\).

So far, we have only considered the case where the integration domain is \( I^s \), or what amounts to the same (modulo a simple change of variable), where the integration domain is an interval. We shall now study a much wider class of integration domains, namely that of bounded Jordan-measurable sets (= sets having a finite elementary volume). By applying, if necessary, a translation and a contraction, we can assume that such an integration domain is contained in \( I^s \). We have already seen in (1.3) what the Monte Carlo approximation should look like in this case. We also noted earlier that Jordan-measurable subsets of \( I^s \) form the most general category of integration domains for which a quasi-Monte Carlo technique can be applied successfully.

The integration error can again be estimated effectively in terms of a suitable notion of discrepancy. We need some preliminaries before we can enunciate this result. Let \( d(\cdot, \cdot) \) be the standard Euclidean distance in \( \mathbb{R}^s \). For a subset \( E \) of \( I^s \) and \( \varepsilon > 0 \), we define

\[
E_\varepsilon = \{ x \in I^s : d(x, y) < \varepsilon \text{ for some } y \in E \},
\]

\[
E_{-\varepsilon} = \{ x \in I^s : d(x, y) > \varepsilon \text{ for all } y \in I^s \setminus E \}.
\]

The set \( E_\varepsilon \) is a superset of \( E \) open in the relative topology in \( I^s \), whereas \( E_{-\varepsilon} \) is a closed subset of \( E \). We classify now the Jordan-measurable subsets \( E \) of \( I^s \) according to the measure of what may be called the “\( \varepsilon \)-collars” of \( E \).

2.10. DEFINITION. Let \( b = b(\varepsilon) \) be a positive nondecreasing function defined for all \( \varepsilon > 0 \) and satisfying \( \lim_{\varepsilon \to 0} b(\varepsilon) = 0 \). Then \( \mathcal{M}_b \) is defined as
the family of all Lebesgue-measurable subsets $E$ of $I^s$ for which the inequalities

$$|E_\varepsilon \setminus E| \leq b(\varepsilon) \quad \text{and} \quad |E \setminus E_{-\varepsilon}| \leq b(\varepsilon)$$

hold for all $\varepsilon > 0$.

It is easily seen that every $E \in \mathcal{M}_b$ is actually Jordan-measurable and that, conversely, every Jordan-measurable subset of $I^s$ belongs to some $\mathcal{M}_b$ (cf. [220, pp. 168–169]). Therefore, the families $\mathcal{M}_b$ provide a complete classification of all Jordan-measurable subsets of $I^s$. The larger the function $b$, the more irregularly shaped can be the sets that are allowed in $\mathcal{M}_b$.

For an integration domain $E \in \mathcal{M}_b$, an error bound for quasi-Monte Carlo integration can now be given in terms of a discrepancy $D_N(\mathcal{M}_c)$ defined according to Definition 2.5, where $c$ is a function closely related to $b$.

2.11. Theorem (Niederreiter [220]). For a function $f$ of bounded variation $V(f)$ on $I^s$ in the sense of Hardy and Krause, an integration domain $E \in \mathcal{M}_b$, and points $x_1, \ldots, x_N$ in $I^s$ we have

$$\left| \frac{1}{N} \sum_{n=1}^{N} f(x_n) - \int_{E} f(t) \, dt \right| \leq (V(f) + |f(1, \ldots, 1)|) D_N(\mathcal{M}_c), \quad (2.10)$$

where $c$ is the function $c(\varepsilon) = b(\varepsilon) + 2\varepsilon$ for $\varepsilon > 0$.

One might wonder why the term $|f(1, \ldots, 1)|$ appears in this estimate. The reason is that, in contrast to earlier inequalities, the left-hand side of (2.10) is not invariant under the shift from $f$ to $f + C$, where $C$ is a constant. Therefore, the upper bound in (2.10) depends also on the magnitude of $f$.

There are two important facts buttressing the usefulness of the families $\mathcal{M}_b$. First, the discrepancy $D_N(\mathcal{M}_b)$ can be estimated effectively in terms of the discrepancy $D_N$, as we shall see in Theorem 3.10, so that we ultimately have control over the size of $D_N(\mathcal{M}_b)$ for the common choices of nodes. Secondly, the discrepancy $D_N(\mathcal{M}_b)$ still satisfies $\lim_{N \to \infty} D_N(\mathcal{M}_b) = 0$ for any infinite sequence uniformly distributed in $I^s$, and uniform conditions on the measure of the “$\varepsilon$-collars” of sets $E$ in a family $\mathcal{M}$ (such as those imposed on the members of $\mathcal{M}_b$) are necessary to guarantee this property for $D_N(\mathcal{M})$ by a result of Billingsley and Topsøe [21]. Thus, the approach based on the families $\mathcal{M}_b$ is essentially the most general one in quasi-Monte Carlo integration.

The case of a convex integration domain deserves special attention since it occurs quite frequently. Information about this class of integration domains is, of course, implicitly contained in the above considerations since every convex subset of $I^s$ belongs to $\mathcal{M}_b$ with $b(\varepsilon) = 2\varepsilon$ for $\varepsilon > 0$. By using special properties of convex sets, it is, however, possible to obtain stronger results.

2.12. Definition. The isotropic discrepancy $J_N$ of the $N$ points $x_1, \ldots, x_N$ in $I^s$ is defined by $J_N = D_N(\mathcal{C})$, where $\mathcal{C}$ is the family of all convex subsets of $I^s$.

The use of $J_N$ was suggested by Hlawka [115] and the term “isotropic discrepancy” was coined by Zaremba [370]. This type of discrepancy figures
2.13. Theorem (Zaremba [370]). If $f$ is a function of bounded variation $V(f)$ on $I^*$ in the sense of Hardy and Krause, $E$ is a convex subset of $I^*$, and $x_1, \ldots, x_N$ are points in $I^*$ with isotropic discrepancy $J_N$, then

$$\left| \frac{1}{N} \sum_{n=1}^{N} f(x_n) - \int_{E} f(t) \, dt \right| \leq (V(f) + |f(1, \ldots, 1)|)J_N.$$ 

Up to now, all the integrands considered were bounded functions. Sobol' [311] succeeded in constructing quasi-Monte Carlo methods for improper integrals in which the behavior of the integrand at the singularity is under control. In a sense, the paper of Hardy and Littlewood [108] may be considered the first contribution to this subject.*

The study of quasi-Monte Carlo integration on the infinite-dimensional unit cube was initiated by Čencov [33] and Sobol' [293]. See also [279], [301], [303, Chapter 7], [313], [314].

The results appearing in this section as well as in the three following ones should be compared with the features of classical techniques for computing simple and multiple integrals. We refer to [53], [95], [173], [329] for expository accounts of these techniques. Summarizing in a nutshell what transpires from this comparison, one can say that conventional methods are without doubt superior in the one-dimensional case, whereas quasi-Monte Carlo methods should be considered more and more favorable as the dimension increases. For simple integrals with integrands of a low degree of regularity, the results are basically equivalent.

Quasi-Monte Carlo methods have been employed for other purposes besides numerical integration. We mention here only those works that are based on the general principles of such methods. In later sections, we will indicate further applications using the special techniques to be described.

There are, of course, several important problems in numerical analysis that can be reduced to the approximate calculation of integrals, and under such circumstances quasi-Monte Carlo methods may be applied readily. A case in point is the numerical solution of integral equations; see [114], [116], [124], [313]. Applications to integro-differential equations occur in [210], [212], to linear partial differential equations in [211], and to the theory of turbulence in [15]. A method of determining surface areas by means of uniformly distributed sequences is discussed in [175]. Interesting connections with complex analysis appear in [118], [123], regarding problems of interpolation and analytic continuation, respectively. Quasi-Monte Carlo methods can also be used to approximate the extreme values of a function ([3], [122], [232], [317]) and the greatest and smallest real part of the eigenvalues of a matrix [318]. On the basis of ideas in [34], applications to the simulation of Markov chains are studied in [312]–[314].

3. Quasi-random points. A scrutiny of the various error estimates in § provides convincing evidence that what we should employ in a quasi-Monte Carlo integration is a set of nodes with very small discrepancy. Finite or
infinite sequences with this property are called \textit{low-discrepancy sequences} and their terms are loosely referred to as \textit{quasi-random points} (or \textit{quasi-random numbers} in the one-dimensional case). Strictly speaking, we have dealt not only with one, but with several concepts of discrepancy, so that the expression \textquote{low-discrepancy sequence} would seem to call for further explanation. It turns out, however, that the various notions of discrepancy are related to $D_N$ in a known manner. Thus, $D_N^*$ and $D_N$ are linked by the inequalities

$$D_N^* \leq D_N \leq 2^s D_N^*$$

in the $s$-dimensional case (see [174, p. 93]), and $J_N$ as well as $D_N(\mathbb{R}_N)$ can also be estimated in terms of $D_N$ (see (3.21) and Theorem 3.10). Moreover, the $L^2$ discrepancy $T_N$ of any $N$ points in $I^s$ satisfies

$$C_s D_N^{(s+2)/2} \leq T_N \leq D_N^*$$

with a constant $C_s > 0$ only depending on $s$ (cf. [220, Theorem 4.2]). Therefore, \textquote{low discrepancy} will be interpreted to mean low discrepancy $D_N$.

The discrepancy of the average sequence of points in $I^s$ is under control because of the law of the iterated logarithm established by Chung [40] in the one-dimensional case and Kiefer [150] in the multidimensional case, according to which we have

$$\limsup_{N \to \infty} \frac{\sqrt{2N}}{\log \log N} D_N^* = 1$$

with probability 1, in the sense of an appropriate product measure on the sequence space.\textsuperscript{10} Together with (3.1) we get

$$D_N = O\left(N^{-1/2} (\log \log N)^{1/2}\right)$$

almost surely. In combination with the error estimates in §2 (in particular, Theorem 2.3 and (2.7)), this ties in rather nicely with the probabilistic error bound for the Monte Carlo method. It is to be expected, however, that clever constructions should produce sequences that behave much better than the average sequence. As a matter of fact, we shall exhibit examples of finite sequences of $N$ points in $I^s$ for which $D_N = O\left(N^{-1}(\log N)^{-1}\right)$. Quasi-Monte Carlo integration with such quasi-random points involves then effective error bounds that are considerably smaller than the probabilistic Monte Carlo bound $O\left(N^{-1/2}\right)$. Thus, as far as numerical integration is concerned, quasi-Monte Carlo methods based on determinate low-discrepancy sequences are, on the whole, preferable to statistical Monte Carlo methods.

We discuss now in detail the relevant properties of the various notions of discrepancy. We start with the one-dimensional case in which all the basic questions have been settled. Here the discrepancy $D_N^*$ can be calculated in an easy manner by arranging the numbers $x_1, \ldots, x_N$ in nondecreasing order of magnitude, which obviously does not affect the value of $D_N^*$. We obtain then the formula\textsuperscript{11}

\textsuperscript{10} See also [30], [242], [243], [372].

\textsuperscript{11} See [215]. For generalizations of this formula referring to other distribution functions or other summation methods, see [219], [222].
\[ D_N^* = \max_{1 \leq n \leq N} \max \left( \left| x_n - \frac{n}{N} \right|, \left| x_n - \frac{n-1}{N} \right| \right) \]
\[ = \frac{1}{2N} + \max_{1 \leq n \leq N} \left| x_n - \frac{2n-1}{2N} \right|. \quad (3.2) \]

We deduce from (3.2) that \( D_N^* > \frac{1}{(2N)} \) for any \( N \) numbers in \( I \), and that \( D_N^* = \frac{1}{(2N)} \) exactly for the sequence

\[ \frac{1}{2N}, \frac{3}{2N}, \ldots, \frac{2N-1}{2N} \]

or one of its rearrangements. A simple argument shows that \( D_N > \frac{1}{N} \) for any \( N \) numbers in \( I \) (cf. [174, p. 90]), and \( D_N = \frac{1}{N} \) holds for the sequence (3.3) and its rearrangements, but also for some other sequences. A formula similar to (3.2) can be obtained for the \( L^2 \) discrepancy \( T_N^2 \), namely

\[ T_N^2 = \frac{1}{12N^2} + \frac{1}{N} \sum_{n=1}^{N} \left( x_n - \frac{2n-1}{2N} \right)^2 \]

provided that \( x_1 < x_2 < \cdots < x_N \) (cf. [220, p. 135]). This shows that the \( L^2 \) discrepancy of any \( N \) numbers in \( I \) is \( > \frac{1}{(\sqrt{12} \cdot N)} \) and that the lower bound is attained precisely for the sequence (3.3) and its rearrangements.\(^{12}\) It is interesting to note that the nodes in (3.3) are also used in a classical integration method, namely the \( N \)-point midpoint rule (see [53, p. 40]). The integration error implied by the use of sequences equal or close to (3.3) was studied by Chui [37], [38], [39].

The terms of the optimal sequence (3.3) depend on the chosen value of \( N \). In computational practice, it is often convenient to be able to change the value of \( N \) without losing the previously obtained data. For this purpose, it is advantageous to work with an infinite sequence and then to take its first \( N \) terms whenever the value of \( N \) has been decided upon. In this way, we may increase \( N \) if we desire greater accuracy and still use the results of the earlier computation.

It requires a deeper analysis to determine the behavior of the discrepancy for infinite sequences. In the first place, the discrepancy \( D_N \) of an infinite sequence cannot be uniformly of the order of magnitude \( N^{-1} \). An important theorem of W. M. Schmidt [282] has established that for any infinite sequence of numbers in \( I \) there exist infinitely many \( N \) such that \( D_N > D_N^* > (\log N)/(100N) \). On the other hand, infinite sequences with \( D_N = O(N^{-1} \log N) \) have been known for several decades. For instance, if \( \alpha \) is an irrational number for which the partial quotients in the continued fraction expansion are uniformly bounded, then the discrepancy \( D_N \) of the sequence \( \{\alpha\}, \{2\alpha\}, \ldots, \{n\alpha\}, \ldots \) of fractional parts\(^{13}\) satisfies \( D_N < CN^{-1}(1 + \log N) \) for all \( N > 1 \), where \( C \) is an explicit constant (compare with [220, Theorem 3.2], [174, p. 125]).

\(^{12}\) A similar statement holds for the \( L^1 \) discrepancy according to a result in [292] that also appeared in an essentially equivalent form in [236, Chapter 10].

\(^{13}\) The fractional part \( \{t\} \) of \( t \in \mathbb{R} \) is defined by \( \{t\} = t - [t] \), where \([t]\) is the greatest integer \( \leq t \). We always have \( 0 \leq \{t\} < 1 \).
Another example, the so-called van der Corput sequence, is more important for two reasons: first, the known discrepancy estimate is better than for the above example, and second, this sequence consists only of dyadic fractions and is therefore perfectly well adapted for the use in binary computers. The sequence can be conveniently defined in terms of a "radical-inverse function". If \( g \geq 2 \) is an integer, then every nonnegative integer \( n \) has an expansion in the base \( g \) of the form

\[
n = \sum_{i=0}^{k} a_i g^i \quad \text{with} \quad a_i \in \{0, 1, \ldots, g - 1\} \quad \text{for} \quad 0 \leq i \leq k,
\]

(3.4)

and this representation is unique apart from adding on higher powers of \( g \) with zero coefficients. Now the radical-inverse function \( \phi_g \) is given by the well-defined expression

\[
\phi_g(n) = \sum_{i=0}^{k} a_i g^{-i-1}.
\]

(3.5)

The action of this function may be described as follows. Write the expansion of \( n \) in the base \( g \) as a string of digits \( a_k \ldots a_1 a_0 \); then \( \phi_g(n) \) results from this by reflection about the "decimal point", i.e., \( \phi_g(n) \) is the \( g \)-adic fraction given by \( 0.a_k a_{k-1} \ldots a_0 \). The condition on the \( a_i \) in (3.4) implies that \( 0 < \phi_g(n) < 1 \) for all \( n > 0 \).

The van der Corput sequence, which made its first appearance in [341], is now defined as the sequence \( \phi_2(0), \phi_2(1), \ldots, \phi_2(n), \ldots \) in \([0, 1)\). According to unpublished results of Tijdeman, its discrepancy \( D_N^* \) satisfies

\[
ND_N^* < \frac{1}{3} \log_2 N + 1 \quad \text{for all} \quad N > 1
\]

and

\[
\lim \sup_{N \to \infty} \left( ND_N^* - \frac{1}{3} \log_2 N \right) \geq \frac{4}{9} + \frac{1}{3} \log_2 3,
\]

where \( \log_2 \) denotes the logarithm to the base 2. The fact that the coefficient \( \frac{1}{3} \) of \( \log_2 N \) is best possible was already shown earlier by Haber [92]. For a quick proof of \( D_N = O(N^{-1} \log N) \), see [174, p. 127]. At present, we know of no infinite sequence in \( I \) whose discrepancy is uniformly smaller than that of the van der Corput sequence.*

There are various ways of estimating the discrepancy of a one-dimensional sequence (see [217] for a survey). The most common procedure is to reduce the estimation of the discrepancy to the problem of estimating associated exponential sums. The theoretical basis for this technique is a general principle of quantitative Fourier inversion. We need the following notion: if \( F \) is a function of bounded variation on \( I \), then its Fourier-Stieltjes transform \( \hat{F} \) is given by

\[
\hat{F}(h) = \int_{0}^{1} e^{2\pi i ht} \, dF(t) \quad \text{for all integers} \quad h.
\]
3.1. **Theorem (Niederreiter and Philipp [234])**. Let $F$ be a nondecreasing function on $I$ with $F(0) = 0$ and $F(1) = 1$, and suppose the function $G$ on $I$ satisfies the Lipschitz condition $|G(u) - G(v)| < L|u - v|$ for $u, v \in I$, as well as $G(0) = 0$ and $G(1) = 1$. Then for every positive integer $m$ we have
\[
\sup_{u, v \in I} \left| (F(v) - F(u)) - (G(v) - G(u)) \right| \leq \frac{4L}{m + 1} + \frac{4}{\pi} \sum_{h=1}^{m} \left( \frac{1}{h} - \frac{1}{m + 1} \right) |\hat{F}(h) - \hat{G}(h)|.
\]

3.2. **Corollary**. Let $x_1, \ldots, x_N$ be any real numbers. Then the discrepancy $D_N$ of the finite sequence of fractional parts $\{x_1\}, \ldots, \{x_N\}$ satisfies
\[
D_N \leq \frac{4}{m + 1} + \frac{4}{\pi} \sum_{h=1}^{m} \left( \frac{1}{h} - \frac{1}{m + 1} \right) \left| \frac{1}{N} \sum_{n=1}^{N} e^{2\pi i h x_n} \right|
\]
for every positive integer $m$.

The corollary results from Theorem 3.1 by setting $F(t) = A([0, t); N)/N$ and $G(t) = t$ for $0 < t < 1$. The inequality in Corollary 3.2, without the specified constants, is due to Erdös and Turán [64]. For a special class of sequences, there exists a simplified version of this inequality which will be useful later on. It is convenient here to introduce for an integer $m > 2$ the summation symbol $\sum_{h \equiv \text{mod } m}$ which designates a sum over the complete residue system mod $m$ consisting of all integers $h$ with $-m/2 < h < m/2$. The summation symbol $\sum_{h \equiv \text{mod } m}$ refers to the same sum, but with $h = 0$ deleted from the range of summation.

3.3. **Lemma**. Let $m > 2$ and $y_0, \ldots, y_{N-1}$ be integers. Then the discrepancy $D_N$ of the finite sequence of fractional parts $\{y_0/m\}, \ldots, \{y_{N-1}/m\}$ satisfies
\[
D_N \leq \frac{1}{m} + \sum_{h \equiv \text{mod } m}^{\ast} \frac{1}{m \sin(\pi|h|/m)} \left| \frac{1}{N} \sum_{n=0}^{N-1} e^{2\pi i h y_n/m} \right|.
\]

To make the relationship with the earlier inequality even more transparent, we use $\sin \pi t > 2t$ for $0 < t < \frac{1}{2}$ and get
\[
D_N \leq \frac{1}{m} + \sum_{h=1}^{[m/2]} \frac{1}{h} \left| \frac{1}{N} \sum_{n=0}^{N-1} e^{2\pi i h y_n/m} \right|.
\]

It should not come as a surprise that we obtain smaller constants in the special case.

The proof of Lemma 3.3 nicely illustrates the principles involved in these types of inequalities. For an integer $k$, let $A(k; N)$ be the number of $n$, $0 < n < N - 1$, with $y_n \equiv k \pmod{m}$. Then we can write
\[
A(k; N) = \sum_{n=0}^{N-1} c_k(y_n),
\]
where $c_k$ is the characteristic function of the coset $k + m\mathbb{Z}$ of $\mathbb{Z}/m\mathbb{Z}$. Now with the abbreviation $e(t) = e^{2\pi it}$ for $t \in \mathbb{R}$ we have
\[ c_k(y) = \frac{1}{m} \sum_{h \equiv 0 \pmod{m}} e(h(y - k)/m) \text{ for } y \in \mathbb{Z}, \]

so that

\[ A(k; N) = \frac{1}{m} \sum_{h \equiv 0 \pmod{m}} e\left(-\frac{hk}{m}\right) \sum_{n=0}^{N-1} e\left(\frac{ny}{m}\right) \]

and

\[ A(k; N) - \frac{N}{m} = \frac{1}{m} \sum_{h \equiv 0 \pmod{m}}^{\star} e\left(-\frac{hk}{m}\right) \sum_{n=0}^{N-1} e\left(\frac{ny}{m}\right). \tag{3.6} \]

Let \( J = [u, v) \) be an arbitrary half-open subinterval of \( I \). We choose the largest closed subinterval of \( J \) of the form \([a/m, b/m]\) with integers \( a < b \), which we denote again by \([a/m, b/m]\). The case where no such subinterval exists can be dealt with easily, since we have then \( A(J; N) = 0 \) and \( v - u < 1/m \), hence

\[ |A(J; N)/N - |J|| = |J| < 1/m. \tag{3.7} \]

In the remaining case, we obtain

\[ A(J; N) - N|J| = \sum_{k=a}^{b} \left( A(k; N) - \frac{N}{m} \right) + \frac{N}{m} (b - a + 1) - N|J| \]

by using (3.6). It follows that

\[ \left| \frac{A(J; N)}{N} - |J| \right| < \frac{1}{m} \sum_{h \equiv 0 \pmod{m}}^{\star} \left| \sum_{k=a}^{b} e\left(hk/m\right) \right| \left| \sum_{n=0}^{N-1} e\left(\frac{ny}{m}\right) \right| \]

\[ + \left| \frac{b - a + 1}{m} - |J| \right|. \tag{3.8} \]

Now for \( 0 < |h| < m/2 \) we have

\[ \left| \sum_{k=a}^{b} e\left(hk/m\right) \right| = \frac{|e(h(b - a + 1)/m) - 1|}{|e(h/m) - 1|} \]

\[ < \frac{2}{|e(h/m) - 1|} = \frac{1}{\sin(\pi|h|m)}. \tag{3.9} \]

From the definition of \( a \) and \( b \) it follows that

\[ a/m = u + \theta_1 \text{ with } 0 < \theta_1 < 1/m \]

and

\[ b/m = v - \theta_2 \text{ with } 0 < \theta_2 < 1/m, \]

so that

\[ \left| \frac{b - a + 1}{m} - |J| \right| = \left| \frac{1}{m} - \theta_1 - \theta_2 \right| < \frac{1}{m}. \]
By combining this with (3.8) and (3.9), we arrive at

\[
\left| \frac{A(J;N)}{N} - |J| \right| < \frac{1}{m} + \sum_{h \pmod{m}}^* \frac{1}{m \sin(\pi |h|/m)} \left| \frac{1}{N} \sum_{n=0}^{N-1} e(hy_n/m) \right|.
\]

In view of (3.7), this inequality holds for all \( J \), and forming the supremum over \( J \) on the left-hand side completes the proof of Lemma 3.3.

There is another inequality for \( D_N \) in terms of exponential sums that is more of theoretical interest. The following general inequality improves and corrects a result of Elliott [63].

3.4. Theorem (Niederreiter [222]). Let \( F \) and \( G \) be two functions on \( I \) satisfying the conditions in Theorem 3.1. Then,

\[
\sup_{u,v \in I} \left| (F(v) - F(u)) - (G(v) - G(u)) \right| < \left( \frac{6L}{\pi^2} \sum_{h=1}^{\infty} \frac{1}{h^2} \left| \hat{F}(h) - \hat{G}(h) \right|^2 \right)^{1/3}.
\]

3.5. Corollary (LeVeque [177]). Let \( x_1, \ldots, x_N \) be any real numbers. Then the discrepancy \( D_N \) of the finite sequence of fractional parts \( \{x_1\}, \ldots, \{x_N\} \) satisfies

\[
D_N < \left( \frac{6}{\pi^2} \sum_{h=1}^{\infty} \frac{1}{h^2} \left| \frac{1}{N} \sum_{n=1}^{N} e^{2\pi i h x_n} \right|^2 \right)^{1/3}.
\]

The difficulties in the multidimensional case start already with the fact that there is no analogue of the simple formula (3.2). One can only provide certain estimates which allow the approximate calculation of \( D_N^* \) in relatively few steps and with a reasonable degree of accuracy (cf. [215]). As a consequence, it is impossible to give an easy description of the finite sequences of \( N \) points in \( I^s \) with minimal discrepancy \( D_N^* \), as was done in (3.3) for \( s = 1 \). Finding these sequences is a laborious computational task, which is nontrivial even for the limited range considered in [361], namely \( s = 2 \) and \( 1 < N < 6 \).

With increasing dimension, the requirement on a sequence to have a small discrepancy becomes more stringent. This is reflected in the known lower bounds for the discrepancy, which depend on the dimension. Roth [261] established a general lower bound on the \( L^2 \) discrepancy \( T_N \), and consequently on \( D_N^* \). In detail, for any \( N \) points in \( I^s \) we have

\[
D_N^* \geq T_N > C_s N^{-1} (\log N)^{(s-1)/2}
\]

with a positive constant \( C_s \) only depending on \( s \). According to [174, p. 102] we may take \( C_s = 2^{-4s}((s - 1) \log 2)^{(1-s)/2} \) for \( s \geq 2 \). An improvement on the order of magnitude of this bound for \( D_N^* \) is only known for \( s = 2 \), where W. M. Schmidt [282] showed

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14 The constant appearing here is better than in the original paper and was obtained in [174, p. 109].
$D^*_N > \frac{\log N}{(132 \log 4) \cdot N}$  \hspace{1cm} (3.11)

for any $N$ points in $I^2$. For infinite sequences of points in $I^s$, it follows from (3.10) and a standard argument (cf. [174, p. 105]) that

$$D^*_N > C'_s N^{-1} (\log N)^{s/2}$$  \hspace{1cm} (3.12)

for infinitely many $N$, where $C'_s$ is an effective positive constant only depending on $s$. A sharper bound is only available for $s = 1$ and was mentioned in the discussion of the one-dimensional case.

The lower bound (3.11) for $s = 2$ is best possible as far as the order of magnitude is concerned. A sequence of $N$ points in $I^2$ with $D^*_N = O(N^{-1} \log N)$ was, for instance, constructed by Roth [261]. It consists of the points $(n/N, \phi_2(n))$, $n = 0, 1, \ldots, N - 1$, where $\phi_2$ is the radical-inverse function from (3.5). In case $N$ is a power of 2, this sequence is known as the Roth sequence. For such $N$, the constant implied in Roth’s estimate was improved by Gabai ([75], [76]) and an exact formula was obtained by Halton and Zaremba [103], in which the leading term for $ND^*_N$ is $(1/3)\log_2 N$ (see [92] for an estimate with this leading term).

In the 2-dimensional case, the lower bound for the $L^2$ discrepancy $T_N$ in (3.10) is also best possible apart from the value of the constant. Curiously enough, this lower bound is not reached by the Roth sequence, for Vilenkin [346] gave an exact formula\(^{15}\) for $NT_N$ with leading term $(1/8)\log_2 N$ (see [92] for an estimate with this leading term). Finite sequences with $T_N = O(N^{-1}(\log N)^{1/2})$ were first constructed by Davenport [50]. His sequences are basically of the form $(n/N, \{n\alpha\})$, $n = 0, 1, \ldots, N - 1$, where $\alpha$ is an irrational with bounded partial quotients in its continued fraction expansion. Examples of such sequences involving only dyadic fractions (as is the case in the Roth sequence) were first obtained by Vilenkin [346] by introducing a perturbation in the Roth sequence. His exact formula for $NT_N$ has leading term $((1/24)\log_2 N)^{1/2}$. Another example (called the Zaremba sequence) yielding the even smaller leading term $((5/192)\log_2 N)^{1/2}$ was found by Halton and Zaremba [103], using a different perturbation of the Roth sequence.\(^ {16}\) The Zaremba sequence may be described as follows. Let $N$ be a power of 2, say $N = 2^k$, $k \geq 1$; then the sequence consists of the $N$ points

$$\left( \frac{a_k}{2} + \frac{a_{k-1}}{2^2} + \cdots + \frac{a_1}{2^k}, \frac{a'_1}{2} + \frac{a'_2}{2^2} + \cdots + \frac{a'_k}{2^k} \right) \in I^2,$$

where $a_1, \ldots, a_k$ take independently of each other the values 0 and 1, while $a'_i = a_i$ for $i$ even, $a'_i = 1 - a_i$ for $i$ odd. Not surprisingly, the discrepancy $D^*_N$ of the Zaremba sequence is smaller than that of the Roth sequence; in fact, the exact formula in [103] for $ND^*_N$ involves the leading term $(1/5)\log_2 N$. See also [100] for a summary of these results. White [360] considered analogues of the Roth and Zaremba sequences using an arbitrary base in the construction (rather than the base 2 in the original sequences) and gave formulas for their

\(^{15}\) This formula was shown independently by Halton and Zaremba [103].

\(^{16}\) See also Roth [262] and Vilenkin [347].
$L^2$ discrepancies (compare also with [357]).

The construction techniques described above can be used for an arbitrary dimension $s > 2$. The so-called Hammersley sequence (of order $N$) was introduced in [105] (see also [99]) and is composed of the $N$ points $(n/N, \varphi_1(n), \ldots, \varphi_{s-1}(n))$, $n = 0, 1, \ldots, N - 1$, in $I^s$, where the bases $g_1, \ldots, g_{s-1}$ of the radical-inverse functions are supposed to be pairwise relatively prime. Usually, one takes for $g_1, \ldots, g_{s-1}$ the first $s - 1$ primes. The discrepancy $D_N$ of the Hammersley sequence was estimated in [99], with the result that $D_N < CN^{-1}(\log N)^{s-1}$ for all $N > 2$, where the constant $C$ only depends on $g_1, \ldots, g_{s-1}$. An infinite sequence with uniformly small discrepancy was constructed in [99] and is nowadays called the Halton sequence. Its terms are the points $(\varphi_1(n), \ldots, \varphi_{s-1}(n))$, $n = 0, 1, \ldots, \text{in } I^s$, where $g_1, \ldots, g_s$ are pairwise relatively prime bases. A convenient algorithm for fast computer generation of these points is available in [102]. The discrepancy $D_N$ of the Halton sequence has the order of magnitude $N^{-1}(\log N)^{s-1}$, as shown in [99]. Better constants were obtained by Meijer [195], and his estimate yields
\[ \frac{N}{D_N} < 2 \sum_{j=1}^s \frac{g_j + (2 \log N)^s}{\log g_j} \text{ for } N \geq 1. \]

A higher-dimensional analogue of the Zaremba sequence was proposed in [357]. We may summarize the above results as follows.

3.6. THEOREM (HALTON [99]). For any dimension $s > 1$, there exists an infinite sequence of points in $I^s$ such that
\[ D_N = O \left( N^{-1} (\log N)^{s-1} \right). \] (3.13)

In addition, for every $N > 2$ there exists a finite sequence of $N$ points in $I^s$ such that
\[ D_N = O \left( N^{-1} (\log N)^{s-1} \right). \] (3.14)

This result is significant since it guarantees that for any dimension there exist quasi-Monte Carlo techniques that perform substantially better than the statistical Monte Carlo method.

It is a widely held belief that the orders of magnitude in (3.13) and (3.14) are best possible. In (3.13) this is known only for $s = 1$ (see the result of W. M. Schmidt [282] quoted in the discussion of the one-dimensional case), and in (3.14) only for $s = 1$ (cf. [174, p. 90]) and $s = 2$ (by (3.11)). Otherwise, there remains the gap between (3.13) and (3.12) on the one hand and between (3.14) and (3.10) on the other.

K. F. Roth has informed the author about an important development concerning the $L^2$ discrepancy $T_N$ for higher dimensions. In his forthcoming paper [263], Roth shows for any $N > 2$ the existence of a sequence of $N$ points in $I^3$ with $T_N = O(N^{-1} \log N)$, and in [264] he establishes for $s > 4$ and any $N > 2$ that there are sequences of $N$ points in $I^s$ with $T_N = O(N^{-1}(\log N)^{(s-1)/2})$. Therefore, the lower bound for $T_N$ in (3.10) is best possible. Calculations of $L^2$ discrepancies for interesting sequences in dimensions $2 < s < 9$ were carried out in [357]. As to the $L^1$ discrepancy,
Sobol' [299] has given a lower bound for any dimension that is of the order $N^{-1}$. This is best possible for $s = 1$. For $s > 1$, W. M. Schmidt [284a] recently obtained an improved lower bound of the order 

$$N^{-1} (\log \log N)/\log \log \log N,$$

and in the same paper it is shown that in (3.10) the $L^2$ discrepancy $T_N$ can be replaced by the $L^p$ discrepancy for any $p > 1$.

The distribution behavior of sequences such as the van der Corput and Roth sequences has led I. M. Sobol' to the development of a general theory of sequences possessing certain uniformity properties with respect to dyadic boxes. Here it is convenient to count also $I$ as a dyadic interval, so that $I$ may now be used in the formation of dyadic boxes (compare with §2). We consider first the case of a finite sequence.\footnote{The Russian literature uses the concise term “net”.}

3.7. DEFINITION (Sobol' [298], [300]). For a nonnegative integer $\tau$, a finite sequence of points in $I^s$ is called a $P_\tau$-net if it contains $2^k$ points, where $k > \tau$ is an integer, and if every dyadic box of volume $2^{s-k}$ contains exactly $2^\tau$ points of the sequence.

For instance, every initial segment of the van der Corput sequence containing $2^k$ points, where $k > 1$, is a $P_0$-net in $I$, and the Roth sequence is a $P_0$-net in $I^2$. The parameter $\tau$ cannot be selected quite arbitrarily since its choice is limited in a certain way by the dimension $s$. For example, there exist arbitrarily long $P_0$-nets in $I$, $I^2$, and $I^3$, but no $P_0$-net in $I^s$, $s > 4$, containing more than 2 points (Sobol' [300]). Let the least value of $\tau$ for which there exist arbitrarily long $P_\tau$-nets in $I^s$ be denoted by $\tau(s)$. Thus, $\tau(1) = \tau(2) = \tau(3) = 0$; moreover, $\tau(4) = 1$ and, in general, $\tau(s) = O(s \log s)$ according to [303, Chapter 6]. The number $\tau(s)$ may be thought of as a geometric characteristic of the $s$-dimensional unit cube. Explicit constructions of $P_\tau$-nets in $I^s$ for $\tau > \tau(s)$ can be found in [300], [303, Chapter 6]. These constructions involve only very simple digital operations with dyadic fractions, so that $P_\tau$-nets may be generated easily in binary computers.

If $\tau$ is small, one would expect a $P_\tau$-net to display an excellent distribution behavior. This is reflected in a general inequality for the nonuniformity $\varphi_\infty(2^k)$ of a $P_\tau$-net in $I^s$ containing $2^k$ points, namely

$$\varphi_\infty(2^k) < 2^{\tau+1+s}.$$

(3.15)

It should be emphasized that this bound does not depend on $k$. In other words, the nonuniformity of a $P_\tau$-net is $O(1)$ in terms of the length of the net. In the light of the error estimate (2.9), this guarantees the usefulness of $P_\tau$-nets for the purposes of numerical integration. We can have equality in (3.15); e.g., if $s = 1, 2, 3$ and $k > s - 1$, then every $P_0$-net in $I^s$ containing $2^k$ points satisfies $\varphi_\infty(2^k) = 2^{s-1}$ (cf. [303, Chapter 6]). There is also a discrepancy estimate for $P_\tau$-nets, to the effect that the discrepancy $D_N^s$ of a $P_\tau$-net in $I^s$ containing $N = 2^k$ points satisfies

$$ND_N^s < 2^\tau \sum_{j=0}^{s-1} \left( \binom{k-\tau}{j} \right)$$

(3.16)
provided that \( k > s - 1 + \tau \) ([298], [300]). Since \( k \) is of the order of magnitude \( \log N \), the inequality (3.16) leads to the estimate \( D_N^k = O((N^{-1}(\log N)^{-1}) \). In other words, the performance of \( P_\tau \)-nets with small \( \tau \) can be compared to that of the Hammersley sequence. In fact, if we take for a given dimension \( s \) a \( P_\tau \)-net in \( I^s \) with minimal \( \tau \), i.e., with \( \tau = \tau(s) \), then an analysis of the constants implied by (3.16) and by the known discrepancy estimate for the Hammersley sequence reveals that, at least for sufficiently large \( s \), such \( P_\tau \)-nets actually "beat" the Hammersley sequence ([300], [303, Chapter 6]). The expected integration error implied by the use of \( P_\tau \)-nets and an integrand chosen at random from a class of functions with rapidly convergent Haar series was estimated in [308], with an improvement in [310].

We turn now to the analogous theory for infinite sequences. We use the following terminology: if \( x_0, x_1, \ldots \) is an infinite sequence, then by a binary segment thereof we mean a block consisting of those \( x_n \) with \( h2^i < n < (h + 1)2^i \), where \( h > 0 \) and \( i > 1 \) are fixed integers.

3.8. Definition (Sobol' [298], [300]). For a nonnegative integer \( \tau \), an infinite sequence of points in \( I^s \) is called an \( L_{P_\tau} \)-sequence if each of its binary segments containing at least \( 2^{\tau+1} \) points is a \( P_\tau \)-net.

The simplest example is provided by the van der Corput sequence, which forms an \( L_{P_0} \)-sequence in \( I \). It is clear from the definitions that \( L_{P_\tau} \)-sequences can only exist for \( \tau > \tau(s) \). But more care has to be taken since we have, for instance, \( \tau(3) = 0 \), but there are no \( L_{P_0} \)-sequences in \( I^3 \); on the other hand, \( I^2 \) still contains \( L_{P_0} \)-sequences (cf. [303, Chapter 6]). As to the explicit construction of \( L_{P_\tau} \)-sequences, the principles pertaining to \( P_\tau \)-nets also apply here. For every dimension \( s \), concrete values of \( \tau \) are known for which \( L_{P_\tau} \)-sequences in \( I^s \) exist.

On the basis of their definition, \( L_{P_\tau} \)-sequences should be very evenly distributed. More explicitly, every \( L_{P_\tau} \)-sequence in \( I^s \) satisfies

\[
\varphi_{\infty}(N) \leq 2^{s-1+\tau} \quad \text{for all } N > 1, \tag{3.17}
\]

so that \( \varphi_{\infty}(N) = O(1) \) as a function of \( N \) (cf. [298], [300]). From the criterion for uniform distribution in terms of \( \varphi_{\infty} \) it follows then that every \( L_{P_\tau} \)-sequence is uniformly distributed. We also infer from (3.17) that the van der Corput sequence satisfies \( \varphi_{\infty}(N) = 1 \) for all \( N > 1 \), a fact already established in [294], and since we have \( \varphi_{\infty}(N) > 2 \) for any \( N \) points in \( I^2 \) with \( N > 2 \) ([303, Chapter 4]), there is a sequence in \( I^2 \) with \( \varphi_{\infty}(N) = 2 \) for all \( N > 2 \). Analogous questions for \( I^3 \) have been discussed in [305]. For the discrepancy \( D_N^\infty \) of an \( L_{P_\tau} \)-sequence in \( I^s \) one has the estimate

\[
ND_N^\infty \leq 2^s \sum_{j=0}^{s-1} \left( m - \tau + 1 \right) j + 1 + 2^s - 1 \quad \text{for } N > 2^{s-1+\tau}, \tag{3.18}
\]

where \( m = \lfloor \log_2 N \rfloor \) (cf. [298], [300]). In terms of orders of magnitude we obtain \( D_N^\infty = O(N^{-1}(\log N)^{\gamma}) \), as for the Halton sequence. If one chooses for a given dimension \( s \) the minimal \( \tau \) for which an \( L_{P_\tau} \)-sequence in \( I^s \) is known to exist, then it turns out that the bound in (3.18) is asymptotically better than the corresponding one for the Halton sequence, in the sense that the implied constant is smaller for sufficiently large \( s \) (cf. [300], [303, Chapter 6]).
**QUASI-MONTE CARLO METHODS**

LP\(_{r}\)-sequences in \(I^s\) possessing additional uniformity properties, e.g., with respect to the \(2^s\) cubes into which \(I^s\) is divided by the hyperplanes \(x_j = 1/2\) (1 \(< j \leq s\)), have been considered in [315], [316].

A program for calculating LP\(_{r}\)-sequences was written up in [317]. Tables for generating such sequences are available in [303, Chapter 6], [309, Appendix], [316] and extend now to the dimension \(s = 51\). The only detailed expository account of the theory of \(P_r\)-nets and LP\(_{r}\)-sequences is in Russian (cf. [303]).

Quasi-Monte Carlo methods based on LP\(_{r}\)-sequences have been employed in various contexts. We mention applications to search algorithms ([3], [309, Chapter 8], [317]), to the problem of bounding the real parts of the eigenvalues of a matrix [318], and to the simulation of systems [353].

Following this discussion of special sequences, we add some remarks concerning the estimation of discrepancy in the multidimensional case. Because of later applications and the fact that the points in many of the interesting finite sequences have rational coordinates, we state a generalization of Lemma 3.3 that is proved in the same way as before. For an integer \(m > 2\), the summation symbol \(\sum_{h \in \mathbb{Z}^s} (\mod m)\) designates a sum over all lattice points \(h = (h_1, \ldots, h_s) \in \mathbb{Z}^s\) with \(-m/2 < h_j < m/2\) for 1 \(< j \leq s\) and \(h \neq 0\). For a lattice point in this range, we define

\[
 r(h, m) = r(h_1, m) \cdots r(h_s, m),
\]

where \(r(h, m) = m \sin(\pi|h|/m)\) for \(h \neq 0\) and \(r(0, m) = 1\).

**3.9. LEMMA.** Let \(y_0, \ldots, y_{N-1}\) be \(s\)-dimensional lattice points. Then for any integer \(m > 2\), the discrepancy \(D_N\) of the finite sequence of fractional parts \(\{1/m\}y_0, \ldots, \{1/m\}y_{N-1}\) satisfies

\[
 D_N \leq \frac{s}{m} + \sum_{h \in \mathbb{Z}^s} \frac{1}{r(h, m)} \left| \frac{1}{N} \sum_{n=0}^{N-1} e^{2\pi i h \cdot y_n} / m \right|,
\]

where the dot denotes the standard inner product in \(\mathbb{R}^s\).

There is also a multidimensional version of the general inequality stated in Theorem 3.1 (see [234]). The special case contained in Corollary 3.2 had been generalized earlier ([157], [331]); see also [174, p. 116]. Extensions of the results in Theorem 3.4 and Corollary 3.5 to higher dimensions are not known.

The isotropic discrepancy \(J_N\) introduced in Definition 2.12 has received some attention recently. A global lower bound was established by W. M. Schmidt [284], namely

\[
 J_N \geq C_s N^{-2/(s+1)}
\]

for any \(N\) points in \(I^s\), where \(C_s\) is a positive constant only depending on \(s\). This is a considerable improvement on a result of Zaremba [370]. The expected order of magnitude of \(J_N\) for a sequence of points selected at random from \(I^s\) is also known. For \(s = 2\), Philipp [244] proved a law of the iterated logarithm for the isotropic discrepancy, and so, in particular, \(J_N = O(N^{-1/2}(\log \log N)^{1/2})\) almost surely. Stute [330] investigated the behavior

\[\text{For } t = (t_1, \ldots, t_s) \in \mathbb{R}^s, \text{ we define } \{t\} = \{(t_1), \ldots, (t_s)\} \in I^s \text{ to be its fractional part.}\]
of \(J_N\) for higher dimensions and obtained the following metric theorems. For \(s = 3\), we have \(J_N = O(N^{-1/2}(\log N)^{3/2})\) almost surely, and for \(s > 4\) we have
\[
J_N = O\left(N^{-2/(s+1)}(\log N)^{2/(s+1)}\right) \quad \text{almost surely.}
\]
These results indicate that the lower bound (3.20) is remarkably accurate. In fact, for \(s > 3\) the exponent of \(N\) in (3.20) cannot be improved, and there is a very thin margin between the performance of sequences with minimal isotropic discrepancy and that of random sequences. The information currently available does not allow us to decide whether (3.20) is also best possible for \(s = 2\). There is also an open question concerning the law of the iterated logarithm for \(J_N\). Clearly, the lower bound (3.20) rules out such a law for \(s > 4\), but it could conceivably hold for \(s = 3\).

At present, one knows of no feasible method of estimating the isotropic discrepancy of interesting sequences of points in \(I^s\), \(s > 2\), in a direct manner. There is, however, a way of getting upper bounds for \(J_N\) via the discrepancy \(D_N\). The best result in this direction is an inequality of Niederreiter and Wills [235] to the effect that
\[
J_N \leq 4s \ D_N^{1/s} \tag{3.21}
\]
for any \(N\) points in \(I^s\). This improves earlier estimates in [111], [119], [206], [217]. An example of Zaremba [367] shows that the exponent \(1/s\) in (3.21) is best possible. On the basis of a theorem of Niederreiter [215], Zaremba [376] presents an algorithm for calculating \(J_N\) in case \(s = 2\). Computational experience with this algorithm is still lacking.

If \(\mathcal{M}_b\) is the family of subsets of \(I^s\) introduced in Definition 2.10, then the discrepancy \(D_N(\mathcal{M}_b)\) defined in accordance with Definition 2.5 can still be estimated in terms of \(D_N\). We recall that this type of discrepancy occurs in the error bound for quasi-Monte Carlo integration over a Jordan-measurable subset of \(I^s\) (see Theorem 2.11). The following general result holds.

3.10. THEOREM (NIEDERREITER AND WILLS [235]). Let \(b = b(\epsilon)\) be a nondecreasing function on the positive reals with \(b(\epsilon) \geq \epsilon\) for all \(\epsilon > 0\) and \(\lim_{\epsilon \to 0^+} b(\epsilon) = 0\). Then for any \(N\) points in \(I^s\) the discrepancies \(D_N(\mathcal{M}_b)\) and \(D_N\) satisfy
\[
D_N(\mathcal{M}_b) \leq 4b(2\sqrt{s} \ D_N^{1/s}).
\]
A somewhat better, but more complicated estimate, which holds also without the condition \(b(\epsilon) \geq \epsilon\), is contained in the original paper. The first inequality of this type was shown in [220]. In the interesting case where \(b(\epsilon) = Ce\) for \(\epsilon > 0\), with \(C > 0\) being a constant, the inequality in [235] yields
\[
D_N(\mathcal{M}_b) \leq (4C\sqrt{s} + 2C + 1)D_N^{1/s}.
\]
For this special case, W. M. Schmidt [284] shows that a similar estimate holds with \(D_N\) replaced by the "cube discrepancy", i.e., the discrepancy \(D_N(\mathcal{W})\) with \(\mathcal{W}\) being the family of closed cubes in \(I^s\) with sides parallel to the coordinate axes.\footnote{W. M. Schmidt also studied the discrepancy with respect to other families of special sets such as triangles, balls, spherical caps, etc. We refer to [283] for an account of this work.}
For certain purposes it has become necessary to construct sequences of
quasi-random points in the infinite-dimensional unit cube $I^\infty$, the cartesian
product of denumerably many copies of $I$. Sobol' [293] defined an infinite-
dimensional version of the Halton sequence, and the same author [301]
introduced $L_p$-sequences in $I^\infty$. Both types of sequences are not only useful
for numerical integration in $I^\infty$ ([293], [301], [303, Chapter 7], [309, Chapter
7], [313], [314]), but also for solving integral equations [313], simulating
Markov chains ([312]–[314]), and for problems arising in particle physics
[302].

4. Good lattice points. The method of low-discrepancy sequences suffers
from a drawback that we have not yet brought into focus, but which is quite
apparent upon inspection of the error bounds in §2. The point is that once the
integrand is sufficiently regular, say of bounded variation in the sense of
Hardy and Krause, then any additional regularity of the integrand is not
reflected in the error estimate. This is in marked contrast to classical
integration methods which, as a rule, become more efficient the more regular
the integrand is. It should be mentioned, though, that the said deficiency is, in
principle, also shared by the statistical Monte Carlo method.

There is a special kind of quasi-Monte Carlo technique, the so-called
method of good lattice points, in which the degree of regularity of the integrand
is duly honored. But a price has to be paid for this luxury since this method
only applies to integrands that are periodic of period 1 in each variable. If $f$
is such an integrand defined on $R^s$, then the approximation used in this method
is of the form

$$
\int_{I^s} f(t) \, dt \approx \frac{1}{m} \sum_{n=1}^{m} f\left(\frac{n}{m} \cdot g\right),
$$

(4.1)

where $m > 2$ is a fixed (large) integer and $g \in Z^s$ is a suitably chosen
$s$-dimensional lattice point. Although the nodes $(n/m)g$ in (4.1) need not
belong to $I^s$, the approximation (4.1) is really of the same type as those
considered earlier since each node may be replaced by its fractional part
$(n/m)g \in I^s$ because of the periodicity property of $f$.

At first glance, the restriction to periodic integrands seems to render the
method unfit for most practical applications. However, there are ways of
transforming a given integrand into a periodic one while preserving regularity
properties and the value of the integral, though this has to be done at the
expense of further calculations. A simple device is the replacement of a
given function $f$ on $I^s$ by the function

$$
h(t_1, \ldots, t_s) = 2^{-s} \sum_{\varepsilon_i=0}^{1} \cdots \sum_{\varepsilon_s=0}^{1} f(\varepsilon_1 + (-1)^{\varepsilon_1} t_1, \ldots, \varepsilon_s + (-1)^{\varepsilon_s} t_s)
$$

for $(t_1, \ldots, t_s) \in I^s$. Then $h$ is identical on opposite faces of $I^s$ and can
therefore be extended to a periodic function on $R^s$ of period 1 in each
variable. Furthermore, we have

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20 This symmetrization process may be seen as an analogue of the antithetic variate technique
in the statistical Monte Carlo method (cf. [28, Chapter 2, §2]).
More sophisticated methods depend on appropriate changes of variables ([109], [169], [275]) or on adding an expression involving Bernoulli polynomials to the integrand ([167], [169], [369]). Surveys of these methods can be found in [369], [373].

The lattice point $g \in \mathbb{Z}^t$ should be chosen so as to yield a small error in (4.1). Let $f$ be a function on $\mathbb{R}^t$ which is periodic with period 1 in each variable and can be expanded into an absolutely convergent multiple Fourier series

$$f(t) = \sum_h c_h e^{2\pi i h \cdot t},$$

where the sum is extended over all lattice points $h \in \mathbb{Z}^t$ and $h \cdot t$ denotes the standard inner product in $\mathbb{R}^t$. If we write again $e(t) = e^{2\pi i t}$ for $t \in \mathbb{R}$ and observe that $c_0 = \langle f \rangle$ then we get

$$\frac{1}{m} \sum_{n=1}^{m} f\left(\frac{n}{m} \cdot g\right) - \int_{I^t} f(t) \, dt = \frac{1}{m} \sum_{n=1}^{m} \sum_h c_h e\left(\frac{n}{m} \cdot h \cdot g\right) - c_0$$

$$= \frac{1}{m} \sum_h c_h \sum_{n=1}^{m} e\left(\frac{n}{m} \cdot h \cdot g\right) - c_0$$

$$= \frac{1}{m} \sum_{h \neq 0} c_h \sum_{n=1}^{m} e\left(\frac{n}{m} \cdot h \cdot g\right).$$

Now the inner sum is equal to 0 if $h \cdot g \equiv 0 \pmod{m}$ and equal to $m$ if $h \cdot g \equiv 0 \pmod{m}$, and so

$$\frac{1}{m} \sum_{n=1}^{m} f\left(\frac{n}{m} \cdot g\right) - \int_{I^t} f(t) \, dt = \sum_{h \neq 0 : h \cdot g \equiv 0 \pmod{m}} c_h. \quad (4.2)$$

Thus the integration error depends on the magnitude of certain Fourier coefficients. As in one-dimensional harmonic analysis, there is a principle in operation which says that the more regular $f$ is, the more quickly its Fourier coefficients $c_h$ will tend to 0 as $h$ moves away from the origin. We measure the distance of $h = (h_1, \ldots, h_t) \in \mathbb{Z}^t$ from the origin by the expression

$$r(h) = \prod_{j=1}^{s} \max(1, |h_j|). \quad (4.3)$$

In other words, $r(h)$ is the absolute value of the product of the nonzero coordinates of $h$.

4.1. DEFINITION. For real numbers $k > 1$ and $C > 0$, we say that $f \in \mathcal{E}_k(C)$ if

$$|c_h| < C r(h)^{-k} \quad \text{for all } h \neq 0, \quad (4.4)$$

and that $f \in \mathcal{E}_k$ if $f \in \mathcal{E}_k(C)$ for some $C > 0$, i.e., if (4.4) holds for some $C > 0$.

There is a sufficient condition for the validity of (4.4) that is easier to check. Let $k > 1$ be an integer and suppose all the partial derivatives
exist and are of bounded variation on \( I^s \) in the sense of Hardy and Krause; then \( f \in \mathcal{S}^k (C) \) with a value of \( C \) which can be given explicitly (Zaremba [366]). Thus (4.4) may be thought of as a regularity condition on \( f \).

If \( f \in \mathcal{S}^k \), then it is seen easily that its Fourier series is automatically absolutely convergent. By combining (4.2) and (4.4), we arrive at the following result which will be basic for the sequel. To facilitate the writing, we set

\[
P^{(k)}(g, m) = \sum_{h \cdot g \equiv 0 \pmod{m}} r(h)^{-k}.
\]

4.2. Theorem (Korobov [161], Hlawka [113]). For every \( f \in \mathcal{S}^k (C) \), \( g \in \mathbb{Z}^s \), and integer \( m > 2 \), we have

\[
\left| \frac{1}{m} \sum_{n=1}^{m} f \left( \frac{n}{m} \right) g - \int_{I^s} f(t) \, dt \right| \leq C P^{(k)}(g, m).
\]

It is clear that in the sum \( P^{(k)}(g, m) \) the main contribution comes from the lattice points \( h \) close to 0. We introduce a simplified, and finite, sum taking into account only such lattice points. We make use of the summation symbol

\[
\sum^*_{h \pmod{m}}
\]

which designates (as in Lemma 3.9) a sum over all \( h = (h_1, \ldots, h_s) \in \mathbb{Z}^s \) with \(-m/2 < h_j < m/2 \) for \( 1 \leq j \leq s \) and \( h \neq 0 \).

4.3. Definition. For \( g \in \mathbb{Z}^s \) and an integer \( m > 2 \), we set

\[
R(g, m) = \sum^*_{h \pmod{m}} r(h)^{-1}.
\]

The sum \( P^{(k)}(g, m) \) can now be estimated in terms of \( R(g, m) \). In fact, by using the method in [174, pp. 156–157] one shows that for any \( k > 1 \), \( g \in \mathbb{Z}^s \), and integer \( m > 2 \) we have

\[
P^{(k)}(g, m) \leq (1 + 2 \zeta(k)) \left( m^{-k} + R(g, m)^k \right),
\]

where \( \zeta(k) = \sum_{n=1}^{\infty} n^{-k} \) is the Riemann zeta-function.

Because of Theorem 4.2 and the estimate (4.6), the usefulness of the present method depends on our ability to find lattice points \( g \) for which \( R(g, m) \) is small. The following new result, which was so far only known for primes \( m \) ([113], [373]), guarantees the existence of such lattice points.

4.4. Theorem (Niederreiter [231]). For any integer \( m > 2 \) and every dimension \( s > 2 \), there exist lattice points \( g \in \mathbb{Z}^s \) with coordinates relatively prime to \( m \) and

\[
R(g, m) < m^{-1} \left( \frac{1}{5} + 2 \log m \right)^s.
\]

A lattice point \( g \in \mathbb{Z}^s \) satisfying (4.7) is called a good lattice point modulo \( m \).

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Korobov [161], [163], [169] speaks of “optimal coefficients modulo \( m \)” when referring to the coordinates of good lattice points modulo \( m \). By combining the above results, we obtain that for a good lattice point \( g \) modulo \( m \) the integration error in (4.1) is of the order of magnitude \( O(m^{-k}(\log m)^{ks}) \) for integrands \( f \in \mathcal{C}^k \). Therefore, the error becomes smaller for larger \( k \), i.e., for more regular integrands. Bahvalov [5] has shown that at least for primes \( m \) there exist lattice points \( g \) for which the exponent \( ks \) of \( \log m \) can be replaced by \( k(s-1) \), but it is also known (see Šarygin [276]) that the exponent cannot be reduced to anything less than \( s-1 \).

The expression defining \( R(g, m) \) does not lend itself to a convenient numerical treatment since it is a sum of many small numbers. A quantity that is much easier to calculate is the integer

\[
\rho(g, m) = \min_h r(h) \quad \text{for} \quad g \in \mathbb{Z}^s, \quad s \geq 2,
\]

where \( h \) ranges over the same lattice points as in the summation for \( R(g, m) \), i.e., \( h = (h_1, \ldots, h_s) \in \mathbb{Z}^s \) with \( h \neq 0 \), \( h \cdot g \equiv 0 \pmod{m} \), and \(-m/2 < h_j < m/2 \) for \( 1 < j < s \). It is trivial that

\[
1/\rho(g, m) < R(g, m).
\]

On the other hand, \( R(g, m) \) may also be estimated from above in terms of \( \rho(g, m) \). The following inequality improves upon earlier results of Hlawka [113] and Zaremba [373].

4.5. **Theorem (Niederreiter [229]).** For any integer \( m \geq 2 \), any dimension \( s \geq 2 \), and any lattice point \( g \in \mathbb{Z}^s \) we have

\[
R(g, m) < \frac{(2 \log m)^s + 4(2 \log m)^{s-1}}{(\log 2)^{s-1}\rho(g, m)} + \frac{2^{s+1}(2^{s+1} - 1)}{(p + s - 2)} \left( \frac{p + s - 2}{s - 1} \right),
\]

where \( p = \lceil \log_2 m \rceil \).

It is now clear that, as far as the quantity \( \rho(g, m) \) is concerned, the desirable lattice points are those for which \( \rho(g, m) \) is large. Such lattice points are also labeled “good lattice points modulo \( m \)”. Thus, the characteristic property of a good lattice point \( g \) modulo \( m \) may be informally described by saying that all lattice points \( h \neq 0 \) solving the congruence \( h \cdot g \equiv 0 \pmod{m} \) should be rather far removed from the origin.\[^{21}\] By combining (4.7) and (4.9), we see that \( \rho(g, m) \) can be at least of the order of magnitude \( m/\log^s m \). A direct approach, first used by Hlawka [116] for primes \( m \), yields an even better result.

4.6. **Theorem (Zaremba [375]).** For every dimension \( s \geq 2 \) and every sufficiently large integer \( m \) there exists at least one lattice point \( g = (1, g_2, \ldots, g_s) \in \mathbb{Z}^s \) such that

\[
\rho(g, m) > (s-1)! m/ (2 \log m)^{s-1}.
\]

At least for primes \( m \), this result can be made effective by using the method

\[^{21}\] This could be thought of as a modulo \( m \) analogue of the notion of “badly approximable point” in simultaneous diophantine approximations (compare with §5).
in [229, §4]. On the other end of the scale, we have \( \rho(g, m) < m/2 \) for any \( g \in \mathbb{Z}^s \), \( s > 2 \), and any integer \( m > 2 \). For if \( g = (g_1, \ldots, g_s) \) with \( g_1 \) relatively prime to \( m \), then \( hg_1 + g_2 \equiv 0 \pmod{m} \) for some \( h \in \mathbb{Z} \) with \(-m/2 < h < m/2\), and so \( \rho(g, m) < \max(1, |h|) < m/2 \). Otherwise, there exists a proper divisor \( d \) of \( m \) with \( dg_1 \equiv 0 \pmod{m} \), and then \( \rho(g, m) < |d| < m/2 \).

The quantity \( P^{(k)}(g, m) \) occurring in the error bound in Theorem 4.2 can be estimated directly in terms of \( \rho(g, m) \). First of all, it is clear from the definitions that for \( s > 2 \) we always have

\[
P^{(k)}(g, m) < P^{(2)}(g, m)/\rho(g, m)^{k-2} \quad \text{for } k > 2.
\]

(4.10)

A detailed study of \( P^{(2)}(g, m) \) leads then to the following inequality.

4.7. **Theorem (Zaremba [373]).** For every dimension \( s > 2 \), any integer \( m > 4 \) and any lattice point of the form \( g = (1, g_2, \ldots, g_s) \in \mathbb{Z}^s \) we have

\[
P^{(2)}(g, m) < \frac{8\pi^2}{6\rho(g, m)^2} \left( \frac{1}{(s-1)!} \left( \log_2 \frac{m}{4} \right)^{s-1} + 2^{s-1} \left( \log_2 \frac{m}{2} \right)^{s-2} \right) + \frac{2\pi^2}{3m^2} \left( 1 + \frac{\pi^2}{3} \right)^{s-1}.
\]

In the 2-dimensional case, the size of the constants appearing in the above estimate can be reduced substantially (cf. [377]). Together with (4.10), we find that for \( s > 2 \) the quantity \( P^{(k)}(g, m) \) is of the order of magnitude \( O(\rho(g, m)^{-k}(\log m)^{s-1}) \).

It must be pointed out that the existence theorems for good lattice points (Theorems 4.4 and 4.6) are by no means constructive, as they depend, respectively, on an averaging argument and a technique of eliminating “bad” lattice points. In general, the only way known so far of producing good lattice points is by an extensive search based on tabulations, preferably of the quantity \( \rho(g, m) \). This is a finite search problem since \( g \) only matters \( \pmod{m} \), so that \textit{a priori} there are \( m^s \) candidates in dimension \( s \). Usually one restricts the attention to lattice points of the form \((1, g_2, \ldots, g_s)\), which further limits the number of possibilities. The task of finding good lattice points is complicated by the fact that their coordinates depend strongly on the dimension. Thus, simplistic schemes such as taking a good lattice point \((g_1, \ldots, g_s)\) modulo \( m \) and searching for an integer \( g_{s+1} \) so as to get a good lattice point \((g_1, \ldots, g_s, g_{s+1})\) modulo \( m \) in dimension \( s + 1 \) are, as a rule, doomed to failure (see however [378]). The reason is, of course, that the nature of the solutions \( h \) of a congruence \( h \cdot g \equiv 0 \pmod{m} \) changes if an additional coordinate is introduced. However, “going down” in the dimension is feasible to some extent because of the following principle. Let \( g^{(s)} = (g_1, \ldots, g_s) \in \mathbb{Z}^s \) with \( s > 3 \) and consider the “truncated” point \( g^{(t)} = (g_1, \ldots, g_t) \), where \( 2 < t < s \); then we have

\[
R(g^{(t)}, m) < R(g^{(s)}, m) \quad \text{for any integer } m > 2,
\]

(4.11)

since with any lattice point \( h^{(t)} = (h_1, \ldots, h_t) \in \mathbb{Z}^t \) contributing to the sum defining \( R(g^{(t)}, m) \) the lattice point \( h^{(s)} = (h_1, \ldots, h_t, 0, \ldots, 0) \in \mathbb{Z}^s \)
contributes to the sum defining $R(g^{(s)}, m)$ and $r(h^{(s)}) = r(h^{(s)})$. For a similar reason we have

$$\rho(g^{(i)}, m) > \rho(g^{(e)}, m) \quad \text{for any integer } m \geq 2. \quad (4.12)$$

If $g^{(o)}$ is now a good lattice point modulo $m$, say with $\rho(g^{(o)}, m)$ large (or even as large as possible), then $\rho(g^{(i)}, m)$ is large, which makes $g^{(i)}$ an acceptable lattice point—but there is, of course, no guarantee that $g^{(i)}$ will have retained the optimality property.

From the constructive point of view, the only case that has been dealt with in a satisfactory manner is $s = 2$. For $s > 3$ one has to take recourse to the currently available tables. The most useful table in the literature is the one of Maisonneuve [185] for dimensions $3 \leq s < 10$ and with the order of magnitude of the moduli $m$ being as large as $10^5$. An earlier table of Saltykov [274] is reprinted in [169] and covers certain prime moduli for dimensions $3 < s < 6$ and moduli that are products of two distinct primes for dimensions $3 < s < 10$. A short table for some selected prime moduli and dimensions $s = 3, 4, 5, 6, 8,$ and $10$ was compiled by Haber [96]. Kedem and Zaremba [148] extended the calculation of Maisonneuve for $s = 3$. Further efforts are reported in [147]. Some of these tables are restricted to good lattice points of the form $(1, g, g^2, \ldots, g^{s-1})$ which are easier to search for since there are initially only $m$ candidates. The use of these special lattice points was suggested by Korobov [165], [166], [169]. They also play an important role in a different context to be expounded later (see §11). Integration errors involved in the calculation of typical 4-dimensional integrals by the method of good lattice points were tabulated by Bruslinskaja [26] (see also [169, Chapter 3]), and for 3- and 4-dimensional integrals by Maisonneuve [185].

As was already mentioned, an explicit construction of good lattice points is possible in the 2-dimensional case. It rests on observing and exploiting an intriguing relation with continued fractions. For an integer $m \geq 2$, we consider a lattice point $g = (1, g)$ with $\gcd(g, m) = 1$. Let

$$\frac{g}{m} = [a_0; a_1, a_2, \ldots, a_q] = a_0 + \frac{1}{a_1 + \frac{1}{a_2 + \ldots + \frac{1}{a_q}}}$$

be the expansion of the rational $g/m$ into a finite simple continued fraction, where we assume $a_q = 1$ for the sake of uniqueness. The integer $a_q$ is just the integral part of $g/m$ and is actually not relevant since we may as well suppose that $1 < g < m$. The positive integers $a_1, a_2, \ldots, a_q$ are the partial quotients of $g/m$. Let

$$K(\frac{g}{m}) = \max(a_1, a_2, \ldots, a_q) \quad (4.13)$$

be the largest partial quotient. Then Zaremba [365] has shown that

$$\frac{m}{K + 2} \leq \rho(g, m) \leq \frac{m}{K} \quad \text{with } K = K\left(\frac{g}{m}\right). \quad (4.14)$$

Consequently, we obtain good lattice points modulo $m$ by choosing rationals
g/m with small partial quotients. For somebody familiar with diophantine approximations, this should not come as a surprise since good lattice points were described earlier as analogues of "badly approximable" points, and this diophantine approximation property is governed by the size of partial quotients.\(^{22}\)

One would expect to get "best" lattice points by selecting rationals g/m with \(K(g/m) = 1\). These rationals are produced via the sequence of Fibonacci numbers, defined by \(F_1 = F_2 = 1\) and \(F_n = F_{n-1} + F_{n-2}\) for \(n > 3\). We put then \(g = F_{n-1}\) and \(m = F_n\), where \(n > 3\), and we have \(K(g/m) = 1\) as desired. This choice of parameters goes back to Bahvalov [5] and Hua and Wang [130]. For the lattice point \(g = (1, F_{n-1})\) we have then \(\rho(g, F_n) = F_{n-2}\) according to a result of Zaremba [365]. Therefore,

\[
\rho(g, F_n) = \left(\frac{F_{n-2}}{F_n}\right) F_n \geq \frac{3}{5} F_n \text{ for } n > 5,
\]

and in the same paper of Zaremba it is verified that for any fixed \(n > 5\) the lattice point \(g\) yields the largest possible value of \(\rho(\cdot, F_n)\). The upshot of this is that for infinitely many moduli \(m\) there are lattice points \(g\) such that \(\rho(g, m)\) is of the order of magnitude \(m\). This is better than what is predicted by Theorem 4.6 for \(s = 2\). It follows from (4.10) and Theorem 4.7 that for these lattice points \(g\) we get \(P^{(k)}(g, m) = O(m^{-k}\log m)\) for \(k > 2\), and so the lower bound of Sarygin for the integration error (which we have mentioned earlier) is attained for \(s = 2\). More concretely, Zaremba [377] has shown for \(g = (1, F_{n-1})\) that

\[
P^{(k)}(g, F_n) < 120 \left(\frac{8}{3}\right)^{k-2} \log F_n \frac{F_n}{F_n^k} \text{ for } k > 2 \text{ and } n > 5.
\]

A detailed study of these lattice points \(g\) was undertaken in [371]. Two-dimensional lattice points arising from continued fraction expansions are also discussed in [98].

The question poses itself whether in the 2-dimensional case \(\rho(g, m)\) can reach the order of magnitude \(m\) for all moduli \(m\) and not just for the Fibonacci numbers. Because of (4.14) this is equivalent to the following problem about rational continued fractions. For \(m > 2\) we set

\[
K_m = \min_{\gcd(g, m) = 1} K\left(\frac{g}{m}\right)
\]

and we ask whether \(K_m\) is uniformly bounded. Zaremba [373, p. 76] has put forth a conjecture which amounts to suggesting that \(K_m < 5\) for all \(m > 2\). This would mean that for every integer \(m > 2\) there exists a reduced fraction \(g/m\) for which all partial quotients are bounded by 5. It would also imply that for every modulus \(m\) there is a lattice point \(g\) with \(\rho(g, m) > m/7\). Borosh [23] calculated \(K_m\) for \(2 < m < 10^4\) and verified Zaremba's conjecture for this range. Actually, his data show that in this range we have \(K_m = 5\) only for \(m = 54\) and \(m = 150\) and \(K_m = 4\) for 23 values of \(m\), the largest one being \(m = 6234\). More recent calculations of Borosh and Niederreiter [24]

\(^{22}\) For instance, for irrational numbers it is known that the smaller the partial quotients, the "worse approximable" the number.
yield $K_m < 3$ for $m = 2^a$, $1 < a < 35$, with $K_m = 2$ occurring frequently. There seems to be ample evidence to conjecture that $K_m < 3$ for all sufficiently large $m$ and that the set of $m > 2$ with $K_m = 2$ has positive lower density. The best order of magnitude presently known is $K_m = O(\log m)$, which follows from Theorem 4.6 and (4.14).

The information available for $s = 2$ suggests that in the general case there may exist good lattice points $g$ modulo $m$ (at least for an infinite class of moduli $m$) such that $\rho(g, m)$ is of the order of magnitude $m/(\log m)^{s-2}$. There is some empirical evidence supporting this conjecture, but nothing has been proved in this direction for $s > 3$.

Attempts have been carried out, however, to explicitly construct lattice points $g \in \mathbb{Z}^r$, $s > 3$, for which $\rho(g, m)$ is at least of a tolerable size. Taking as their cue the analogy between good lattice points and badly approximable points (in the sense of simultaneous diophantine approximations), Hua and Wang [134] employ the following procedure. Let $\alpha = (\alpha_2, \ldots, \alpha_s) \in \mathbb{R}^{s-1}$ have coordinates which are real algebraic numbers such that $1, \alpha_2, \ldots, \alpha_s$ are linearly independent over the rationals. For a fixed integer $m > 2$, let $g_2, \ldots, g_s$ be integers such that we have the simultaneous diophantine approximation

$$|\alpha_j - g_j/m| < dm^{-s/(s-1)} \quad \text{for} \quad 2 < j < s,$$

where $d$ is a positive constant. Then it is shown that the lattice point $g = (1, g_2, \ldots, g_s)$ satisfies

$$\rho(g, m) > c(\alpha, e, d, s)m^{s/2(s-1)-\varepsilon}$$

(4.16) for every $\varepsilon > 0$, where $c(\alpha, e, d, s)$ is a positive constant depending on the indicated parameters. As to a convenient choice of $\alpha$, the authors already proposed in [132], [133] to consider, for a prime $2s + 1$, the numbers

$$\alpha_j = 2 \cos \frac{2\pi(j - 1)}{2s + 1} \quad \text{for} \quad 2 < j < s,$$

which together with 1 form a basis for the real cyclotomic field $\mathbb{Q}(2 \cos 2\pi/(2s + 1))$ of degree $s$ over $\mathbb{Q}$. Calculations of Haber [96] used this particular point $\alpha$ and led to promising results.

Another construction of Hua and Wang [134] is based on the properties of a special class of algebraic numbers. Let $\omega$ be a PV number (for Pisot-Vijayaraghavan number) of degree $s > 2$, i.e., $\omega$ is a real algebraic integer of degree $s$ such that $\omega > 1$ and its remaining algebraic conjugates $\omega^{(2)}, \ldots, \omega^{(s)}$ all lie in the open unit disk. We may order these conjugates so that

$$|\omega^{(2)}| < |\omega^{(3)}| < \cdots < |\omega^{(s)}| < 1.$$  

(4.17)

Suppose the minimal polynomial of $\omega$ over $\mathbb{Q}$ is $x^s - a_{s-1}x^{s-1} - \cdots - a_1x - a_0 \in \mathbb{Z}[x]$, and let $b = (b_0, b_1, \ldots, b_{s-1}) \in \mathbb{Z}^s$ be a lattice point $\neq 0$. Define the linear recurring sequence $Q_0, Q_1, \ldots$ of integers by $Q_i = b_i$ for $0 < i < s - 1$ and the recurrence relation

$$Q_{n+s} = a_{s-1}Q_{n+s-1} + \cdots + a_1Q_{n+1} + a_0Q_n \quad \text{for} \quad n = 0, 1, \ldots.$$

Then for sufficiently large $n$ we have $|Q_n| > 2$ and the lattice point $g = (1,
On the other hand, it satisfies
\[ \rho(g, |Q_n|) > c(\omega, b, \epsilon)|Q_n|^{(1+\beta)/2-\epsilon} \]
for every \( \epsilon > 0 \), where \( c(\omega, b, \epsilon) \) is a positive constant depending on the indicated parameters and
\[ \beta = -\log|\omega(\epsilon)|/\log \omega. \]

The stipulation (4.17) implies \( \beta < 1/(s-1) \), and so the lower bound in (4.18) is not better than the one in (4.16). However, the present method has the advantage that it is more explicit than the previous one. It is suggested to take for \( \omega \) the largest real root of the irreducible equation
\[ x^s - x^{s-1} - \cdots - x - 1 = 0 \]
and \( b = (0, \ldots, 0, 1) \). The resulting linear recurring sequence \( Q_0, Q_1, \ldots \) may then be thought of as a generalized Fibonacci sequence. For \( s = 2 \) we obtain the example involving Fibonacci numbers.

Good lattice points can be used to produce finite sequences with small discrepancy. Let \( g \in \mathbb{Z}^s \) be a lattice point and \( m > 2 \) a modulus. It follows then from Lemma 3.9 and the inequality
\[ r(h, m) > 2r(h) \]
for lattice points \( h = (h_1, \ldots, h_s) \neq 0 \) with \(-m/2 < h_j < m/2 \) (1 \( \leq j \leq s \)) that the discrepancy \( D_m \) of the points \( \{(1/m)g\}, \{(2/m)g\}, \ldots, \{(m/m)g\} \) satisfies
\[ D_m < s/m + \frac{1}{2}R(g, m). \]
Thus, if \( g \) is chosen to be a good lattice point modulo \( m \), we get a small discrepancy \( D_m \). Theorem 4.4 implies the following result.

**4.8. Theorem.** For every dimension \( s > 2 \) and every integer \( m > 2 \), there exists a lattice point \( g \in \mathbb{Z}^s \) with coordinates relatively prime to \( m \) such that the discrepancy \( D_m \) of the points \( \{(1/m)g\}, \{(2/m)g\}, \ldots, \{(m/m)g\} \) satisfies
\[ D_m < \frac{s}{m} + \frac{1}{2m} \left( \frac{3}{2} + 2 \log m \right)^2. \]

One conjectures that for certain \( s \)-dimensional lattice points one even has
\[ D_m = O \left( m^{-2}(\log m)^{s-1} \right), \]
which would be the same order of magnitude as that of the discrepancy of the Hammersley sequence (see §3). This conjecture has only been verified for \( s = 2 \), where, predictably, one takes a Fibonacci number \( m = F_n \) (\( n > 3 \)) and the lattice point \( g = (1, F_{n-1}) \) and one obtains
\[ D_m < (7/6)m^{-1}\log(15m) \]
according to a result of Zaremba [365]. Sobol' [291], [294], [296] investigated the behavior of sequences arising from good lattice points with respect to his nonuniformities \( \varphi_q \) and \( \varphi_\infty \) (see also [303, Chapter 5]). It follows from (2.8) that the finite sequence of points in Theorem 4.8 satisfies \( \varphi_\infty(m) = O(\log^2 m) \), a result shown in [294] for primes \( m \). Proinov [249] considers the \( L^2 \) discrepancy of the points obtained from those in Theorem 4.8 after symmetrization with respect to all vertices of \( I^s \).

A number of other questions about good lattice points have been studied in the literature. The effect of randomizing the lattice point \( g \) was considered by Bahvalov [8], [9] and Stojancev [324]. The device of shifting the nodes by a random vector is studied in [49]. Solodov [321] demonstrated the applicability
of the method of good lattice points for certain integration domains with smooth boundary. Further theoretical results and/or remarks about the subject can be found in [7], [11], [42], [97], [128], [162], [165], [229], [269], [319], [320], [322], [367], [374], [383]. For surveys of various aspects of the theory of good lattice points, see [10], [95], [131], [168], [169], [174, Chapter 2, §5], [303, Chapter 5], [356], [368], [373].

The method of good lattice points can be used in several areas of numerical analysis. A rather immediate application is to the approximate solution of integral equations. This was first considered by Korobov [164], with later efforts in this direction being contained in [114], [116], [124], [134], [166], [169, Chapter 4], [265]-[268], [275], [356], [381]. Interpolation techniques based on good lattice points were first studied by Rjaben'kiï [257] and Smoljak [287], and this was further pursued in [117], [166], [169, Chapter 4], [266], [270], [276], [288], [355]. An application of these interpolation formulas is made in [139]. Bahvalov [6] uses good lattice points in a Dirichlet problem; see also [326], [380], [382]. The papers [258] and [325] discuss similar applications to Cauchy problems. A relation between good lattice points and approximation theory occurs in [4]. Korobov [170] employs good lattice points to approximate the values of multiple sums with many terms.

Prior to his investigation of good lattice points, Korobov [160] considered nodes of the form \((n/m, n^2/m, \ldots, n^s/m)\), \(n = 1, 2, \ldots, m\), where \(m = p\) or \(p^2\) with a prime \(p > s\) (see also [78], [113], [116], [169, Chapter 2]). These do not perform as well as good lattice points since for an integrand \(f \in \mathcal{E}^k\), \(k > 1\), we only get an error bound of \(O(m^{-1/2})\). The discrepancy of this finite sequence was estimated by Hlawka [113], [116], with the result that \(D_m = O(m^{-1/2}\log^s m)\).

A set of nodes that is sometimes used in classical integration techniques consists of the points \((n_1/m, n_2/m, \ldots, n_s/m)\), where \(m > 2\) is an integer and the \(n_j\) run independently through all integers from 0 to \(m - 1\), so that there are altogether \(N = m^s\) nodes. Korobov [162] notes that in a quasi-Monte Carlo integration with these points and for an integrand \(f \in \mathcal{E}^k\), \(k > 1\), the integration error is \(O(N^{-k/s})\). See also [113], [116], [129], [169, Chapter 1], [289], [383]. For other sets of nodes related to this one or to good lattice points, we refer to [13], [51], [74], [128], [271], [272].

5. Application of diophantine approximations. In the method of good lattice points we use as nodes the multiples of a "rational" point \((1/m)g \in \mathbb{R}^s\). A closely related method is obtained if one employs instead the multiples of an "irrational" point \(\alpha \in \mathbb{R}^s\), i.e., of a point \(\alpha = (\alpha_1, \ldots, \alpha_s)\) for which 1, \(\alpha_1, \ldots, \alpha_s\) are linearly independent over the rationals. The periodicity conditions on the integrand will be the same as in §4. Thus, the approximation formula can be written as

\[
\int_{I^s} f(t) \, dt \approx \frac{1}{N} \sum_{n=1}^{N} f(n\alpha).
\] (5.1)

In order to analyze the integration error, we assume again that \(f\) belongs to the function class \(\mathcal{E}^k(C)\) for some \(k > 1\) and \(C > 0\) (see Definition 4.1). By expanding \(f\) into its absolutely convergent multiple Fourier series, an argu-
ment similar to the one in the preceding section leads to
\[
\left| \frac{1}{N} \sum_{n=1}^{N} f(n\alpha) - \int_{I^*} f(t) \, dt \right| = \left| \frac{1}{N} \sum_{n=1}^{N} c_n \sum_{h \neq 0} e^{2\pi i h(n\alpha)} \right| \leq \frac{C}{N} \sum_{h \neq 0} r(h)^{-k} \left| \sum_{n=1}^{N} e^{2\pi i h \alpha} \right|.
\]

The inner sum is a finite geometric series with \( h \cdot \alpha \not\in \mathbb{Z} \) because of the condition on \( \alpha \). A straightforward estimate yields then
\[
\left| \sum_{n=1}^{N} e^{2\pi i h \alpha} \right| \leq \frac{1}{2\|h \cdot \alpha\|},
\]
where \( \|t\| = \min_{m \in \mathbb{Z}} |t - m| \) for \( t \in \mathbb{R} \)
designates the absolute distance from \( t \) to the nearest integer. Altogether we have
\[
\left| \frac{1}{N} \sum_{n=1}^{N} f(n\alpha) - \int_{I^*} f(t) \, dt \right| \leq \frac{C}{2N} \sum_{h \neq 0} r(h)^{-k} \|h \cdot \alpha\|^{-1}.
\]

The convergence of this infinite series depends on how small \( \|h \cdot \alpha\| \) can become for lattice points \( h \) close to the origin. The study of this behavior belongs to the realm of simultaneous diophantine approximations. We use the following notion to classify irrational points with regard to their diophantine approximation character.

5.1. **Definition.** For a real number \( \eta \), an irrational point \( \alpha \in \mathbb{R}^s \) is said to be of finite type \( \eta \) if \( \eta \) is the infimum of all numbers \( \sigma \) for which there exists a positive constant \( c = c(\alpha, \sigma) \) such that
\[
r(h)^\sigma \|h \cdot \alpha\| > c
\]
holds for all lattice points \( h \in \mathbb{Z}^s \) with \( h \neq 0 \). If no such number \( \sigma \) exists, then \( \alpha \) is said to be of infinite type.

It follows from the Minkowski linear forms theorem that we always have \( \eta \geq 1 \). Irrational points of the smallest possible type \( \eta = 1 \) may be called badly approximable points. Such points can be constructed explicitly. A fundamental theorem of W. M. Schmidt [281] shows that every algebraic irrational point\(^{23}\) is of finite type \( \eta = 1 \). A result of A. Baker [12] implies that \( \alpha = (e^{r_1}, \ldots, e^{r_s}) \) with distinct nonzero rationals \( r_1, \ldots, r_s \) is of finite type \( \eta = 1 \).

We return now to the estimate (5.2). Using techniques from [216], one proves the convergence of the infinite series on the right-hand side under the assumption that \( k \) is larger than the type of \( \alpha \). Thus, we obtain the following error estimate.

\(^{23}\) This is an irrational point with real algebraic numbers as coordinates.
5.2. Theorem (Hasegrove [109], Niederreiter [220]). Let \( \alpha \in \mathbb{R}^s, s > 1 \), be an irrational point of finite type \( \eta \). Then we have

\[
\frac{1}{N} \sum_{n=1}^{N} f(n\alpha) - \int_{I} f(t) \, dt = O\left( \frac{1}{N} \right)
\]

for every \( f \in \mathcal{E}^k \) with \( k > \eta \).

For greater efficiency one will choose a point \( \alpha \) of the smallest possible type \( \eta = 1 \). The use of algebraic irrational points was proposed in [216] on the basis of the pertaining special case of Theorem 5.2. More specifically, one may, in case \( 2s + 3 \) is prime, take the point

\[
\alpha = \left( 2 \cos \frac{2\pi}{2s + 3}, 2 \cos \frac{4\pi}{2s + 3}, \ldots, 2 \cos \frac{2\pi s}{2s + 3} \right)
\]

arising from a basis of a real cyclotomic field (see §4); for any \( s > 1 \) one may choose \( \alpha = (\xi, \xi^2, \ldots, \xi^s) \) with \( \xi = p^{1/(s+1)} \) and \( p \) prime, or

\[
\alpha = (\sqrt{p_1}, \ldots, \sqrt{p_s})
\]

with distinct primes \( p_1, \ldots, p_s \). The first two possibilities are to be preferred since the coordinates stem from an algebraic number field of degree \( s + 1 \) over \( \mathbb{Q} \) (which is the smallest degree that can occur for an algebraic irrational point), whereas the coordinates of the third point generate an algebraic number field of degree \( 2s \) over \( \mathbb{Q} \). Zinterhof [384] works with the irrational points of A. Baker mentioned above.

Historically, this method dates back to papers of Richtmyer [254] and Peck [240], the latter stating the result of Theorem 5.2 as a conjecture. Similar ideas were used later on in [5], [10], [14], [16], [18], [44], [45], [162], [169, Chapter 2], [197], [358].

Computational work based on this method was carried out in [52] and [259], and for a slightly modified method in [93]. For an application of this method to interpolation problems, see [384].

If we have a very regular integrand, say \( f \in \mathcal{E}^k \) with \( k \) much larger than the type \( \eta \) of \( \alpha \), then the method, as it stands, does not honor the high degree of regularity of \( f \). A way has been found to alleviate this deficiency, namely by abandoning the equal-weight formula (5.1) and working with weight distributions adapted to the specific regularity condition on \( f \). This is the only instance in the theory of quasi-Monte Carlo integration where the use of nonequal weights leads to a demonstrable gain in accuracy. The idea of constructing special weight distributions to achieve better error bounds in results such as Theorem 5.2 occurs already in Bahvalov [5], Haselgrove [109], and Wang [356].

A general theorem along these lines is obtained as follows. For positive integers \( q \) and \( N \), we define weights of the form \( a_{Nn}^{(q)} N^{-q} \) with \( 0 < n < q(N - 1) \). The \( a_{Nn}^{(q)} \) are positive integers which are introduced most conveniently by means of a generating function. In detail, the \( a_{Nn}^{(q)} \) are determined from the polynomial identity

\[
\left( \sum_{j=0}^{N-1} z^j \right)^q = \sum_{n=0}^{q(N-1)} a_{Nn}^{(q)} z^n.
\]

There is also a result of Pjateckiï-Šapiro included in [77] and appearing as well in [169, Chapter 2].
With these weights, the following result holds.

5.3. Theorem (Niederreiter [220]). Let \( q \) be a positive integer and \( \alpha \in \mathbb{R}^s \), \( s > 1 \), an irrational point of finite type \( \eta \). Then we have

\[
\frac{1}{N^q} \sum_{n=0}^{q(N-1)} a_{\alpha}^{(n)} f(n\alpha) - \int_{I^s} f(t) \, dt = O\left( \frac{1}{N^q} \right)
\]

for every \( f \in \mathcal{S}^k \) with \( k > q \eta \).

Thus, while carrying out only about \( q \) times as many function evaluations, the error bound becomes approximately the \( q \)th power of what we had earlier, at least for sufficiently regular integrands. If we take for \( \alpha \) an algebraic irrational point, then the integration error is \( O(N^{-q}) \) for every \( f \in \mathcal{S}^k \) with \( k > q \). A related result was established by Hua and Wang [134].

It is an advantageous feature of this method that the point \( \alpha \) is chosen independently of \( N \). Thus, if the value of \( N \) is increased, the previously calculated function values can be used again. This is not the case in the method of good lattice points, since the property of being a good lattice point modulo \( m \) is obviously relative to the choice of \( m \). Of course, both methods share the drawback that a nonperiodic integrand first has to be “periodized” by one of the procedures mentioned in §4 before the integration technique can be applied. It should also be mentioned that in an actual machine calculation based on the method in this section one will have to replace the irrational \( \alpha \) by a rational approximation. This produces an additional inaccuracy which may, however, be classified as a roundoff error. A further limitation of the method was discussed by Tsuda [339] who used numerical data to show that it does not work so well if the integrand has a very high and narrow peak. The intrinsic reason is that the constant involved in the error term becomes too large in this case.

An analysis of the proofs of Theorems 5.2 and 5.3 reveals that the only noneffective constant is the one stemming from the diophantine inequality (5.3). In the multidimensional case, it is usually very difficult to get an effective value for \( c(\alpha, \sigma) \) as long as \( \sigma \) is close to the type \( \eta \) of \( \alpha \). For algebraic irrational points \( \alpha \), there are ways of obtaining an effective constant \( c(\alpha, \sigma) \) by considering larger values of \( \sigma \) (cf. [216]). A fuller discussion of the issue of effectiveness can be found in [220, §10].

The discrepancy of the sequence \( \{\alpha\}, \{2\alpha\}, \ldots, \{n\alpha\}, \ldots \) of fractional parts can be estimated in terms of the type of \( \alpha \). If \( \alpha \in \mathbb{R}^s \) is an irrational point of finite type \( \eta \), then we have

\[
D_N = O(N^{-1/(\eta+\varepsilon)} + \varepsilon)
\]

for every \( \varepsilon > 0 \) (cf. [220, §6], [174, pp. 131–132]). This is known to be essentially best possible for \( s = 1 \) (cf. [174, p. 124]). If \( \alpha \in \mathbb{R}^s \) is an algebraic irrational point, then the above estimate implies \( D_N = O(N^{-1+\varepsilon}) \) for every \( \varepsilon > 0 \), a result conjectured and heuristically supported by Richtmyer [255] and first shown by Niederreiter [217]. Further remarks on this case can be found in [256]. For \( s = 1 \), we already mentioned in §3 that there are irrationals for which \( D_N = O(N^{-1} \log N) \). In fact, most irrationals will
produce a discrepancy nearly as small as that, for if \( g \) is a positive nondecreasing function such that \( \sum_{n=1}^{\infty} g(n)^{-1} \) converges, then \( D_N = O(N^{-1}(\log N) g(\log \log N)) \) for almost all \( \alpha \in \mathbb{R} \) in the sense of Lebesgue measure (cf. [219], [174, p. 128]). In the multidimensional case, W. M. Schmidt [280] proved a metric result to the effect that for every \( \varepsilon > 0 \) we have \( D_N = O(N^{-1}(\log N)^{s+1+\varepsilon}) \) for almost all \( \alpha \in \mathbb{R}^s \) in the sense of \( s \)-dimensional Lebesgue measure. No individual \( \alpha \in \mathbb{R}^s, s > 2 \), is known for which \( D_N = O(N^{-1}(\log N)^{s+1}) \).

PART II. PSEUDO-RANDOM NUMBERS

6. Random numbers vs. pseudo-random numbers. We have seen in §1 that the essential step in the statistical Monte Carlo method is the random sampling of points (or numbers) from a given set by independent trials. If we assume, for simplicity, that the set is the unit interval \( I = [0, 1] \), then an "unbiased" execution of this sampling procedure should produce a sequence of "random numbers" in \( I \). However, as long as no concrete definitions of "unbiased" and no equal opportunity regulations for numbers are adopted, the above concept of a sequence of random numbers is fictional, or at best an idealization. We are faced here with an obvious contradiction in terms, namely to try to define or characterize a supposedly haphazard process by a fixed set of rules known in advance. This dilemma shifts the problem from the mathematical to the metaphysical level, without resolving it, of course (to the relief of every gambling casino in the world). The question has indeed attracted the attention of philosophers of science ([152], [194], [246, Chapter 6]). Some authors\(^{25}\) have also perceived the challenge of getting a handle on the vexing idea of randomness as being basic for the foundations of probability theory.

Many attempts at "defining" a sequence of random numbers in \( I \) recognize the following principles: (i) the sequence should satisfy certain distribution properties; (ii) these distribution properties should be invariant under certain selection rules for subsequences. This viewpoint is due essentially to von Mises ([350], [351, §1]), although (i) was acknowledged much earlier (see, e.g., [344, pp. 64–67]). On this basis, Knuth [154, §3.5] discusses in detail a hierarchy of proposed definitions for a random number sequence (see also [198]). The minimum requirement is that of uniform distribution in \( I \). But even a very nicely distributed sequence can exhibit a distinctly nonrandom behavior, e.g., the van der Corput sequence (see §3) has the property that its terms alternately hit the intervals \([0, \frac{1}{2})\) and \(\left[\frac{1}{2}, 1\right)\). Therefore, the statistical independence of successive terms has to be taken into account, and this leads eventually to the requirement of complete uniform distribution.\(^{26}\) A sequence \( x_1, x_2, \ldots \) of elements of \( I \) is called completely uniformly distributed (abbreviated CUD) if for every integer \( s > 1 \) the sequence of points \( (x_n, x_{n+1}, \ldots, x_{n+s-1}) \), \( n = 1, 2, \ldots \), is uniformly distributed in \( I^s \). Equivalently, the sequence \( x_1, x_2, \ldots \) is CUD if for every \( s > 1 \) and every \( s \)-dimensional lattice point \( (h_1, \ldots, h_s) \neq 0 \) the sequence of fractional parts \( \{h_1 x_n + \ldots\)
\[ h_2x_{n+1} + \cdots + h_2x_{n+2} \}, \ n = 1, 2, \ldots, \] is uniformly distributed in \( I \).

CUD sequences have, in fact, been used as sources of random numbers for Monte Carlo-type calculations and simulations (see, e.g., [15], [34], [73], [309, Chapter 7]). Explicit constructions of CUD sequences have already been given in the older literature on the subject (cf. [247]), but these are too inconvenient for practical use. A CUD sequence consisting only of dyadic fractions was constructed by Knuth [153], and a principle of obtaining simple sequences that are in a sense approximately CUD was pointed out by Haber [94]. More recently, Rauzy [252] proved that if \( f \) is an entire function that is not a polynomial, attains real values on the real axis, and satisfies

\[
\lim_{r \to \infty} \frac{\log \log M(f; r)}{\log \log r} < \frac{3}{4},
\]

where \( M(f; r) = \sup_{|z| < r} |f(z)| \), then the sequence of fractional parts \( \{f(1)\}, \{f(2)\}, \ldots, \{f(n)\}, \ldots \) is CUD. M. B. Levin [180] showed how to construct, for every transcendental \( \alpha \) such that the sequence of fractional parts \( \{\alpha\theta\}, \{\alpha\theta^2\}, \ldots, \{\alpha\theta^n\}, \ldots \) is CUD. This is related to a metric result of Franklin [72] to the effect that \( \{\theta\}, \{\theta^2\}, \ldots, \{\theta^n\}, \ldots \) is CUD for almost all transcendental numbers \( \theta > 1 \) (in the sense of Lebesgue measure).

So far, we have ignored the invariance principle (ii) mentioned above. If one insists that the sequence and all of its subsequences should be CUD, one arrives at a void concept, since any CUD sequence has a subsequence of terms in \([0, \frac{1}{2}]\), which is therefore not even uniformly distributed in \( I \). Hence, some restriction has to be imposed on the admissible selection rules. If one allows only recursively enumerable selection rules, then the same argument as before leads to the interesting situation that every concretely given sequence fails this test for randomness. Further possibilities for selection rules are discussed in [154, §3.5], and one also finds there A. Wald's construction of a sequence enjoying remarkable invariance properties. A systematic study of the relationship between complete uniform distribution (or "normality", the analogous concept for sequences of digits) and selection rules is carried out in [145], [146].

For finite sequences of digits, an information-theoretic approach to a definition of randomness based on preparatory work of Kolmogorov [159] has been chosen by Chaitin [35] and Martin-Löf [191]. The idea here is to designate those sequences as random which, among all sequences of fixed length, maximize the shortest length of the program required to generate the sequence on a Turing machine.\(^2\) A pleasant introduction to this concept is presented in [36], where one also finds an elaboration on the paradoxical statement that randomness in this sense cannot be verified for a specified sequence of digits. A detailed survey of this and other approaches to randomness is given in [285]. See also the recent papers [178], [179].

\(^2\) It is easily seen that a sequence of the form \( \{\alpha\theta\}, \{\alpha\theta^2\}, \ldots, \{\alpha\theta^n\}, \ldots \) cannot be CUD for algebraic \( \theta \).

\(^2\) As Knuth [154, p. 149] points out, this is of course the worst concept from the viewpoint of practical random-number generation.
There is a way of generating "random numbers" (in the intuitive sense) that is very much in the spirit of the statistical Monte Carlo method, namely by using physical devices. Among the possibilities that have been considered are white noise produced by electronic circuits ([28, Chapter 6], [90]), counts of the emission of radiated particles ([28, Chapter 6], [138], [203], [349]), flipping coins [354], mechanical gadgets operating on the principle of the roulette wheel ([104], [142]), and automated versions of the latter ([149], [251]). Extensive tables recording these experiments or reproducing random numbers obtained from other data have been published ([144], [149], [251], [335]). See [334] for the early history of these tables.

The use of physical or tabulated random numbers is problematic because of the difficulties of storage, the necessity of frequent testing for randomness, and the fact that they are not generated in the computer, but by an external source. We have also seen that those concepts of randomness that sound convincing from an axiomatic standpoint suffer from the practical deficiency that no concrete sequence actually used in a calculation can verifiably satisfy the definition. Therefore, one has taken recourse to sequences that make no pretense of being "random" in any meaningful sense of the term, but which can be readily generated in the computer by simple arithmetic algorithms, while still passing an assortment of statistical tests for randomness. The terms of such sequences are collectively (and loosely) called pseudo-random numbers (abbreviated PRN). It should be emphasized right away that no such sequence can perform well under all imaginable tests for randomness. Rather, the user of PRN has to be aware of the specific statistical properties that are desirable in his Monte Carlo calculation and choose PRN that are known to pass these tests. This "relativity principle" will become amply clear given the results of §11. On the whole, PRN have a record of meeting any reasonably limited set of statistical requirements if adequately chosen for the particular purpose.

We will only discuss PRN simulating the uniform distribution on I. A variety of methods has been devised for transforming such PRN into others obeying a nonuniform distribution law, e.g., a normal distribution. We refer to [2], [28, Chapter 7], [66, Chapter 2], [80], [142], [154, §3.4], [193, Chapter 3], [213]. The bibliographies on pseudo-random number generation mentioned in §1 also cover these techniques.

As we pointed out, it should be a cardinal virtue of a sequence of PRN that it gets a stamp of approval from a collection of statistical tests for randomness. We review now some of the tests that have been considered frequently in the literature. 29 We assume throughout that \(x_0, x_1, \ldots\) is a (finite or infinite) sequence of numbers in I to be tested for "randomness".

A. Equidistribution test. Here we check the evenness of distribution in I of an initial segment \(x_0, x_1, \ldots, x_{N-1}\) of the sequence. This is done by calculating (or estimating) the discrepancy \(D_N\) (see §2) of these numbers. The test is passed if \(D_N\) is small. In an alternative method, sometimes called the frequency test, one partitions I into subintervals, counts the number of terms falling into the various subintervals, and performs a chi-square test on these data.

29 This list is by no means exhaustive. See [154, §3.3] for more information.
B. Gap test. Let \( J \) be a fixed proper subinterval of \( I \). If for some \( n > 0 \) we have \( x_{n+j} \notin J \) for \( 0 < j < k - 1 \), but \( x_{n+k} \in J \), we speak of a gap of length \( k \). We choose a positive integer \( h \) and count the number of gaps of lengths 0, 1, \ldots, \( h - 1 \) and \( > h \) until a large total number of gaps is reached. Then we apply a chi-square test using the probabilities \( p_i = p(1-p)^i \) for \( 0 < i < h - 1 \) and \( p_h = (1-p)^h \) for these respective categories, where \( p \) is the length of \( J \).

C. Run test. A segment of the sequence satisfying \( x_n < x_{n+1} < \cdots < x_{n+k-1} \), but \( x_{n-1} > x_n \) and \( x_{n+k-1} > x_{n+k} \), is called a "run up" of length \( k \). We count the number of "runs up" of lengths 1, 2, \ldots, \( h \) and \( > h + 1 \) in a given initial segment of the sequence. Since adjacent runs are not independent, we cannot use a straightforward chi-square test for these data. A more complicated statistic has to be computed (see [154, pp. 60–63]). A similar test can be performed for "runs down".

D. Permutation test. We choose an integer \( s > 2 \) and consider the \( s \)-tuples \( (x_n, x_{n+1}, \ldots, x_{n+s-1}) \), \( 0 < n < N \). There are \( s! \) possible relative orderings among the entries of such an \( s \)-tuple. We count the number of times each ordering appears and perform a chi-square test, using the probability \( 1/s! \) for each ordering, or determine the maximal deviation from the expected number.

E. Serial correlation. This is a rather weak test for the interdependence between \( x_n \) and \( x_{n+1} \). We calculate the serial correlation coefficient

\[
\sigma_1 = \frac{N \sum_{n=0}^{N-1} x_n x_{n+1} - \left( \sum_{n=0}^{N-1} x_n \right)^2}{N \sum_{n=0}^{N-1} x_n^2 - \left( \sum_{n=0}^{N-1} x_n \right)^2}.
\]

If \( x_n \) and \( x_{n+1} \) are almost independent, then \( |\sigma_1| \) is very small. It is a deficiency of this test that the converse does not necessarily hold. One may, more generally, consider the serial correlation coefficient \( \sigma_h \) reflecting the interdependence between \( x_n \) and \( x_{n+h} \).

F. Spectral test. This test was introduced in [48] (see also [47]) and depends on a Fourier transform technique. We assume that all PRN in the given sequence are rationals with common denominator \( m \), as will be the case for the important method to be introduced in §7. For any \( s \)-dimensional lattice point \( (h_1, \ldots, h_s) \) we define the limit

\[
L(h_1, \ldots, h_s) = \lim_{N \to \infty} \frac{1}{N} \sum_{n=0}^{N-1} e(h_1 x_n + h_2 x_{n+1} + \cdots + h_s x_{n+s-1}),
\]

which will exist, for instance, in the special case mentioned above. In an ideal case of randomness we should have \( L(h_1, \ldots, h_s) = 1 \) if \( h_1 \equiv 0 \cdots \equiv h_s \equiv 0 \pmod{m} \) and \( L(h_1, \ldots, h_s) = 0 \) otherwise. The deviation from this behavior is a measure for the nonrandomness of the given sequence of PRN. The difficulty here is to find a realistic way of putting this idea into a quantitative form. In [48] an intuitive analogy with the theory of harmonic oscillation was used (see also [154, §3.3.4]), although there is no mathematical basis for this procedure.

G. Serial test. This is the most reliable test for statistical independence of
successive terms. For a fixed dimension \( s > 2 \), we consider the points\(^{30}\)
\[ x_n = (x_n, x_{n+1}, \ldots, x_{n+s-1}), \quad n = 0, 1, \ldots, N - 1, \quad \text{in } I^s. \]
For a truly random sequence these points will, in the long run, be very evenly distributed over \( I^s \). Therefore, the quality of the PRN can be measured by the discrepancy \( D_N \) (see §2) of the points \( x_0, \ldots, x_{N-1} \). This test has an added significance since information about \( D_N \) yields effective error bounds for quasi-Monte Carlo integrations using the nodes \( x_0, \ldots, x_{N-1} \) (compare with §2). An older and less powerful version of this test operates like the frequency test mentioned under A, by starting from a partition of \( I^s \) into subintervals and proceeding as in A.

For the special class of PRN that will be introduced in the next section, there is also the possibility of studying the structure of the lattice generated by the points occurring in test G (see §10 for details).

7. Linear congruential pseudo-random numbers. There is conceivably an overwhelming multitude of techniques for generating PRN, some discovered and many still dormant, but a simple algorithm proposed by Lehmer [176] has easily become the most popular one and is now a traditional tool of the numerical analyst. This method can be described as follows. Let \( m > 2 \) and \( r \) be integers, let \( y_0 \) be an integer in the least residue system modulo \( m \), i.e., \( 0 < y_0 < m \), and let \( \lambda \) be a positive integer relatively prime to \( m \) which may as well be taken from the least residue system modulo \( m \). Then a sequence \( y_0, y_1, \ldots \) of integers in the least residue system modulo \( m \) is generated by the recursion
\[ y_{n+1} \equiv \lambda y_n + r \pmod{m} \quad \text{for } n = 0, 1, \ldots \quad (7.1) \]
From this sequence we derive a sequence \( x_0, x_1, \ldots \) of numbers in \( I \) by setting \( x_n = y_n / m \) for \( n = 0, 1, \ldots \), and this is already the desired sequence of so-called linear congruential PRN (abbreviated LCPRN). In this context, \( m \) is referred to as the modulus, \( \lambda \) as the multiplier, and \( r \) as the increment. In practice, \( m \) is taken to be a large prime or a large power of 2, the latter choice being particularly suitable for a binary computer. To rule out trivial cases, we suppose that \( \lambda \not\equiv 1 \pmod{m} \) and \( (\lambda - 1)y_0 + r \not\equiv 0 \pmod{m} \).\(^{31}\)
An explicit formula for \( x_n \) can easily be established by induction, namely
\[ x_n = \left\lfloor \frac{\lambda^n y_0 + \frac{\lambda^n - 1}{\lambda - 1} \cdot \frac{r}{m}}{m} \right\rfloor \quad \text{for } n = 0, 1, \ldots \quad (7.2) \]
It is customary to distinguish the homogeneous case \( r \equiv 0 \pmod{m} \) (also called the multiplicative congruential method) and the inhomogeneous case \( r \not\equiv 0 \pmod{m} \) (also called the mixed congruential method). In the homogeneous case we can assume w.l.o.g. that \( y_0 \) is relatively prime to \( m \), for otherwise the same sequence of LCPRN could be produced with a smaller modulus.

The parameters \( m, \lambda, r, \) and \( y_0 \) have to be chosen in such a way that the resulting PRN will pass appropriate statistical tests for randomness. The

\(^{30}\) In [154, pp. 55-56] only nonoverlapping \( s \)-tuples are considered, but this restriction is not necessary.

\(^{31}\) If \( (\lambda - 1)y_0 + r \equiv 0 \pmod{m} \), then \( y_1 \equiv y_0 \pmod{m} \), and the sequences \( y_0, y_1, \ldots \) and \( x_0, x_1, \ldots \) are constant.
modulus \( m \) is selected in accordance with machine capabilities, typical choices being \( m = 2^{32} \) or \( m = 2^{35} \). The initial value \( y_0 \) has practically no, and the increment \( r \) only little influence on the behavior of the PRN, so that the properties of the sequence are mainly governed by the choice of the multiplier \( \lambda \). In the light of the simplicity of the generation method, it is always a surprise to find out how remarkably well these LCPRN perform under rather stringent tests once a multiplier has been selected intelligently.

We review now briefly the elementary properties of LCPRN. See [135], [142], [154, §3.2] for a more detailed analysis. Most importantly, the sequence \( x_0, x_1, \ldots \) is always purely periodic. We shall use \( \tau \) to denote its least period. In the homogeneous case, it is easily seen that \( \tau \) is equal to the least positive integer \( n \) for which \( \lambda^n \equiv 1 \pmod{m} \), i.e., \( \tau \) is the exponent to which \( \lambda \) belongs modulo \( m \). In the inhomogeneous case and for a prime or prime power modulus, \( \tau \) can be determined quickly on the basis of (7.2). We first introduce some notation that will also be convenient later on. For \( m = p^\alpha, p \) prime, \( \alpha > 1 \), let \( \kappa \) be the largest integer such that \( p^\kappa \) divides \( \lambda - 1 \) and let \( \omega \) be the largest integer such that \( p^\omega \) divides \( (\lambda - 1)y_0 + r \). Then the following result is contained in [226, Lemma 8] if we note that under our conditions we always have \( \omega \leq \kappa \).

7.1. **Lemma.** Let \( m = p^\alpha, p \) prime, \( \alpha > 1 \). Then the period \( \tau \) of a sequence of LCPRN with modulus \( m \) and multiplier \( \lambda \) is equal to the exponent to which \( \lambda \) belongs modulo \( p^{\alpha-\omega+\kappa} \).

If \( m \) has at least two distinct prime factors, the period can still be obtained from this result because of [154, p. 16, Lemma Q]. For the homogeneous case and \( m = p^\alpha \), with \( p = 2, \alpha < 2, \) or \( p \) an odd prime, the largest possible period is \( \tau = p^{\alpha-1}(p - 1) \), and it is attained precisely for multipliers \( \lambda \) that are primitive roots modulo \( m \), whereas for \( m = 2^\alpha, \alpha > 3 \), the largest possible period is \( \tau = 2^{\alpha-2} \), attained precisely for \( \lambda \equiv \pm 3 \pmod{8} \) if \( \alpha > 4 \) and \( \lambda \equiv 3, 5, 7 \pmod{8} \) for \( \alpha = 3 \). For general \( m \), the maximum period is obtained from the quoted result in [154]. In the inhomogeneous case, we can actually have \( \tau = m \), i.e., all fractions in \([0, 1)\) with denominator \( m \) are generated, and this happens precisely if the following requirements are met: (i) \( r \) is relatively prime to \( m \); (ii) \( \lambda \equiv 1 \pmod{p} \) for every prime \( p \) dividing \( m \); (iii) \( \lambda \equiv 1 \pmod{4} \) if 4 divides \( m \). This criterion can be deduced from Lemma 7.1, but it follows also from other elementary considerations (cf. [135], [154, p. 15]).

We have dwelt on the question of the period since a good sequence of LCPRN should at least have a long period. One reason is, of course, that periodicity is not a typical feature of a truly random sequence. Therefore, a sequence with a short period can hardly pretend to be pseudo-random. Moreover, in an actual quasi-Monte Carlo calculation based on LCPRN we want to use a large number \( N \) of terms, but we must have \( N \leq \tau \) so that the grossly nonrandom periodicity property does not come into play; hence \( \tau \) should be large. Further indications why a large value of \( \tau \) is preferable will be given later on (see, e.g., §9).

In interesting cases we may describe all fractions that are generated by the method. This is, of course, trivial whenever the period is \( m \). If \( \lambda \) is a primitive
root modulo $m$ and we are in the homogeneous case, then we produce exactly all reduced fractions in $[0, 1)$ with denominator $m$. Again in the homogeneous case and for $m = 2^n$ with $n > 3$ and $\lambda \equiv 5 \pmod{8}$, we obtain exactly all fractions in $[0, 1)$ of the form $a/m$ with $a \equiv y_0$ or $3y_0 \pmod{8}$. A thorough study is carried out in [188], where it is observed, for instance, that a sequence of LCPRN is always made up of a block of less than $m$ terms, followed by translates of that block. The frequent phenomenon of one sequence of LCPRN being a cyclic permutation of another is investigated in [31].

Historically, Lehmer’s method was not the first proposal for pseudorandom number generation. A few years earlier, J. von Neumann already used his “middle-square method” to produce pseudo-random numbers for Monte Carlo calculations (cf. [154, §3.1], [352]). However, this technique has proved unsatisfactory since it tends to lead to short cycles and the expected distribution is nonuniform (cf. [336]). In more promising schemes that may be called hybrid methods, one starts from a sequence of LCPRN and adds a new twist or combines it with other sequences of PRN. An intriguing idea is the “shuffling” scheme of MacLaren and Marsaglia [184] (see also [189]). Here one takes a sequence of LCPRN and shuffles it in a way which is directed by a different sequence of PRN. Thus we get a sequence of PRN with a “pseudo-random shuffle” thrown in for good measure. Theoretical evidence that shuffling may improve the performance of LCPRN is presented in [260]. Other hybrid methods can be found, for instance, in [32], [82], [88], [202], [273], [290], [297], [359].

Another idea that has been brought up in the literature is to replace (7.1) by a higher-order recurrence relation. Here it is convenient to use a prime number $p$ as a modulus. Then a congruence such as (7.1) can be viewed as a recurrence relation in the finite field $F_p = \mathbb{Z}/p\mathbb{Z}$. Let $k$ be a positive integer which will serve as the order of the new recurrence relation. We note that the finite field $F_{p^k}$ of $p^k$ elements is an extension of $F_p$ and that its multiplicative group $F_{p^k}^*$ is cyclic. A polynomial $f(x) = x^k - a_kx^{k-1} - \cdots - a_0 \in \mathbb{Z}[x]$ is called a primitive polynomial modulo $p$ if the polynomial $\hat{f}(x) \in F_p[x]$ canonically associated with $f(x)$ is the minimal polynomial over $F_p$ of a generator of $F_{p^k}^*$. With such a primitive polynomial modulo $p$, we set up the $k$th order homogeneous linear recurrence

$$y_{n+k} \equiv a_{k-1}y_{n+k-1} + \cdots + a_0y_n \pmod{p} \quad \text{for } n = 0, 1, \ldots .$$

Any sequence $y_0, y_1, \ldots$ of integers in the least residue system modulo $p$ satisfying (7.3) with $(y_0, \ldots, y_{k-1}) \neq (0, \ldots, 0)$ is called a maximal period sequence modulo $p$. The reason for this terminology lies in the fact that the length of the period of a maximal period sequence modulo $p$ is equal to $p^k - 1$, the largest possible period of any $k$th order homogeneous linear recurring sequence in $F_p$. A maximal period sequence modulo $p$ is easily seen to be purely periodic (see [379] for these facts about linear recurring sequences).

From a maximal period sequence $y_0, y_1, \ldots$ modulo $p$ we derive a sequence $x_0, x_1, \ldots$ of PRN in $I$ by setting $x_n = y_n/p$ for $n = 0, 1, \ldots$. In
practice, the prime \( p \) will be chosen very large, but \( p = 2 \) is sometimes used to produce a pseudo-random sequence of binary digits. The advantages offered by such sequences of PRN can be seen as follows. Consider the \( k \)-tuples \((y_n, y_{n+1}, \ldots, y_{n+k-1})\), \( n = 0, 1, \ldots, p^k - 2 \). They are all distinct (for otherwise the period of \( y_0, y_1, \ldots \) would be less than \( p^k - 1 \)) and \((0, \ldots, 0)\) cannot occur among them, hence they run exactly through all \( k \)-tuples \((0, \ldots, 0)\) of elements in the least residue system modulo \( p \). Therefore, the points \((x_n, x_{n+1}, \ldots, x_{n+k-1})\), \( n = 0, 1, \ldots, p^k - 2 \), show an excellent distribution in \((p^k \mathbb{Z}/\mathbb{Z})^k\). Also, in a full period of \( y_0, y_1, \ldots \) each nonzero residue modulo \( p \) occurs exactly \( p^k - 1 \) times and 0 occurs exactly \( p^k - 1 \) times; each pair \((0, 0)\) of residues modulo \( p \) occurs exactly \( p^k - 2 \) times among \((y_n, y_{n+1})\), \( n = 0, 1, \ldots, p^k - 2 \), whereas \((0, 0)\) occurs exactly \( p^k - 2 \) times; and so on for \( s \)-tuples with \( s < k \). Altogether, we can say that the sequence \( x_0, x_1, \ldots \) of PRN leads to practically perfect distribution properties in \((p^k \mathbb{Z}/\mathbb{Z})^k\) for all dimensions \( 1 < s < k \). Of course, we have to pay a price for this benefit, in the sense that the generation method (7.3) is more complicated than (7.1). We will report on further results about these generators in §§9, 10.

The idea of using higher-order recurrences for pseudo-random number generation can be traced back to [62], [81], [343]. The method (7.3) was proposed by Tausworthe [333] and endorsed by Knuth [154, §3.2.2]. We shall refer to the associated sequences of PRN as Tausworthe generators. So-called "additive generators", for which the coefficients in the recurrence are either 0 or 1, emerged already in the early times of pseudo-random number generation. A standard example is the Fibonacci generator \( y_{n+2} \equiv y_{n+1} + y_n \) (mod \( m \)) (cf. [332]), which was, however, soon discarded because it displays a notoriously bad behavior. The recursion \( y_{n+k} \equiv y_{n+k-1} + y_n \) (mod \( m \)) is a decent generator for \( k = 16 \) according to [85].

After Tausworthe's paper, higher-order generators received more attention. Apart from the investigations on Tausworthe generators to be mentioned later, there are studies of higher-order recurrences modulo powers of 2 (cf. [245]) and of certain second-order recurrences modulo primes (cf. [58]). A special method of generating sequences of \( n \)-bit strings of binary digits by a second-order recurrence involving bit-by-bit addition modulo 2 and cyclic shifts was proposed in [250], and algebraic results about the periods of these sequences were shown in [277].

8. Exponential sums. The main tool in our treatment of the equidistribution test and serial test for LCPRN will be certain exponential sums that are intimately connected with these numbers. That such sums should play a rôle is quite clear in the light of Lemmas 3.3 and 3.9. We consider the general case of a \( k \)th order linear recurring sequence \( z_0, z_1, \ldots \) of integers satisfying the recurrence relation

\[
z_{n+k} = a_{k-1}z_{n+k-1} + \cdots + a_0z_n + r \quad \text{for } n = 0, 1, \ldots,
\]

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where \( a_{k-1}, \ldots, a_0, r \) are fixed integers. If such a sequence is viewed modulo an integer \( m \geq 2 \), it looks like a sequence \( y_0, y_1, \ldots \) used in the generation of PRN (see §7). In this case, we will assume, for simplicity, that \( a_0 \) and \( m \) are relatively prime.\(^{34}\) This guarantees that \( z_0, z_1, \ldots \) is purely periodic modulo \( m \). Let \( \tau = \tau(m) \) be the length of its least period modulo \( m \), and let \( \mu = \mu(m) \) be the length of the least period modulo \( m \) of the sequence \( v_0, v_1, \ldots \) of integers satisfying

\[
v_{n+k} = a_{k-1}v_{n+k-1} + \cdots + a_0v_n \quad \text{for } n = 0, 1, \ldots,
\]

with the initial values \( v_0 = \cdots = v_{k-2} = 0, v_{k-1} = 1 \) (\( v_0 = 1 \) if \( k = 1 \)). The following is a special case of a more general result. As before, we write \( e(t) = e^{2\pi it} \) for \( t \in \mathbb{R} \).

8.1. **Theorem (Niederreiter [224], [225]).** Let \( m \geq 2 \) and \( b \) be relatively prime integers, let \( z_0, z_1, \ldots \) be the \( k \)th order linear recurring sequence of integers satisfying (8.1) with \( \gcd(a_0, m) = 1 \), and let \( \tau \) and \( \mu \) be defined as above. Then,

\[
\left| \sum_{n=0}^{\tau-1} e\left( \frac{b}{m} z_n \right) \right| \leq \left( \frac{m^k \tau - \tau^2}{\mu} \right)^{1/2} \quad (8.2)
\]

and

\[
\left| \sum_{n=0}^{N-1} e\left( \frac{b}{m} z_n \right) \right| \leq \left( \frac{m^k \tau}{\mu} \right)^{1/2} \left( \frac{2}{\pi} \log \tau + \frac{2}{5} \right) + \frac{N}{\tau} \left( \frac{m^k \tau - \tau^2}{\mu} \right)^{1/2} \quad \text{for } 1 < N < \tau. \quad (8.3)
\]

The upper bounds in (8.2) and (8.3) represent slight improvements on the original inequalities, which are achieved by subtracting the contribution from the zero vector in [224, Equation (16)] (compare with the proof of Theorem 8.2). These improvements are of interest when \( \tau \) is large.

If a sequence of LCPRN is generated by the multiplicative congruential method, we are led to consider the recurrence relation \( z_{n+1} = \lambda z_n, \ n = 0, 1, \ldots \), which produces powers of \( \lambda \), up to a constant factor. We present a proof for this case to indicate the ideas involved. Note that under the conditions of the subsequent result we have \( \tau = \mu \).

8.2. **Theorem (Korobov [172], Niederreiter [226]).** Let \( m \geq 2, b, \lambda, \) and \( c \) be integers with \( \gcd(b, m) = \gcd(\lambda, m) = 1 \) and \( \lambda \) belonging to the exponent \( \tau \) modulo \( m \). Then,

\[
\left| \sum_{n=0}^{\tau-1} e\left( \frac{b}{m} \lambda^n \right) e\left( \frac{cn}{\tau} \right) \right| \leq \begin{cases} (m - \tau)^{1/2} & \text{if } c \equiv 0 \pmod{\tau}, \\ m^{1/2} & \text{if } c \equiv 0 \pmod{\tau}, \end{cases} \quad (8.4)
\]

and

\[
\left| \sum_{n=0}^{N-1} e\left( \frac{b}{m} \lambda^n \right) \right| \leq m^{1/2} \left( \frac{2}{\pi} \log \tau + \frac{2}{5} \right) + \frac{N}{\tau} (m - \tau)^{1/2} \quad \text{for } 1 < N < \tau. \quad (8.5)
\]

\(^{34}\) This is always satisfied in the applications to PRN.

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PROOF. For an integer \( a \), write
\[
\sigma(a, c) = \sum_{n=0}^{\tau-1} e\left( \frac{a}{m} \lambda^n \right) e\left( \frac{cn}{\tau} \right).
\]
The general term of this sum, considered as a function of \( n \), is periodic with period \( \tau \). Therefore, for any integer \( j \geq 0 \) we have
\[
\sigma(a, c) = \sum_{n=0}^{\tau-1} e\left( \frac{a}{m} \lambda^{n+j} \right) e\left( \frac{c(n+j)}{\tau} \right),
\]
and so we obtain the transformation formula
\[
|\sigma(a, c)| = \left| \sum_{n=0}^{\tau-1} e\left( \frac{a\lambda^j}{m} \lambda^n \right) e\left( \frac{cn}{\tau} \right) \right| = |\sigma(a\lambda^j, c)|. \quad (8.6)
\]
Since the integers \( b, b\lambda, \ldots, b\lambda^{\tau-1} \) are pairwise incongruent modulo \( m \) and not divisible by \( m \), (8.6) implies that
\[
|\sigma(b, c)|^2 = \sum_{j=0}^{\tau-1} |\sigma(b\lambda^j, c)|^2 < \sum_{a=1}^{m-1} |\sigma(a, c)|^2
\]
\[
= \sum_{a=1}^{m-1} |\sigma(a, c)|^2 - |\sigma(0, c)|^2
\]
\[
= \sum_{h,j=0}^{\tau-1} e\left( \frac{c(h-j)}{\tau} \right) \sum_{a=1}^{m-1} e\left( \frac{a}{m} (\lambda^h - \lambda^j) \right) - |\sigma(0, c)|^2
\]
\[
= m\tau - |\sigma(0, c)|^2.
\]
The inequalities in (8.4) are immediate consequences. To prove (8.5), we use the identity
\[
\sum_{n=0}^{N-1} e\left( \frac{b}{m} \lambda^n \right) = \sum_{n=0}^{\tau-1} e\left( \frac{b}{m} \lambda^n \right) \sum_{y=0}^{N-1} e\left( \frac{1}{\tau} \sum_{c=1}^{\tau} e\left( \frac{c(n-y)}{\tau} \right) \right)
\]
for \( 1 \leq N \leq \tau \),
which follows from the fact that the sum over \( y \) is equal to 1 for \( 0 \leq n < N - 1 \) and equal to 0 for \( N \leq n < \tau - 1 \). We rewrite this identity in the form
\[
\sum_{n=0}^{N-1} e\left( \frac{b}{m} \lambda^n \right) = \frac{1}{\tau} \sum_{c=1}^{\tau} \left( \sum_{y=0}^{N-1} e\left( \frac{cy}{\tau} \right) \right) \left( \sum_{n=0}^{\tau-1} e\left( \frac{b}{m} \lambda^n \right) e\left( \frac{cn}{\tau} \right) \right).
\]
Then by (8.4),
\[
\left| \sum_{n=0}^{N-1} e\left( \frac{b}{m} \lambda^n \right) \right| < \frac{1}{\tau} \sum_{c=1}^{\tau} \left| \sum_{y=0}^{N-1} e\left( \frac{cy}{\tau} \right) \right| \left| \sum_{n=0}^{\tau-1} e\left( \frac{b}{m} \lambda^n \right) e\left( \frac{cn}{\tau} \right) \right|
\]
\[
< \frac{1}{\tau} m^{1/2} \sum_{c=1}^{\tau-1} \left| \sum_{y=0}^{\tau-1} e\left( \frac{cy}{\tau} \right) \right| + \frac{N}{\tau}(m-\tau)^{1/2},
\]
and (8.5) is obtained by a straightforward estimation of the remaining sum (compare with [226, Lemma 2]).
The sums in (8.4) contain Gaussian sums to a prime modulus or to an odd prime power modulus as special cases. The second part of (8.4) reduces then to a familiar inequality for Gaussian sums, and for certain values of \( c \) we will even have equality (cf. [225, p. 60]). On the other hand, Gaussian sums can be used to prove Theorem 8.2, but yield nothing for the inhomogeneous case or for higher-order recurrences. As an illustration, we prove the following result on the basis of ideas in [224]. We use \( \phi \) to denote Euler’s totient function.

8.3. THEOREM. Let \( m, b, \) and \( \lambda \) be as in Theorem 8.2. Then,

\[
\left| \sum_{n=0}^{\tau-1} e\left( \frac{b}{m} \lambda^n \right) \right| < \sqrt{m} - \frac{\tau}{\phi(m)} \left( \sqrt{m} - 1 \right) \tag{8.7}
\]

and

\[
\left| \sum_{n=0}^{N-1} e\left( \frac{b}{m} \lambda^n \right) \right| < \sqrt{m} \left( \frac{2}{\tau} \log \tau + \frac{2}{3} \right) + N \left( \frac{\sqrt{m}}{\tau} - \frac{\sqrt{m} - 1}{\phi(m)} \right)
\]

for \( 1 \leq N \leq \tau \). \tag{8.8}

PROOF. We observe that \( e\left( \frac{bh}{m} \right) \), considered as a function on the integers \( h \) with \( \gcd(h, m) = 1 \), can be expanded into a finite Fourier series with respect to the Dirichlet characters \( \psi \) modulo \( m \). Thus,

\[
e\left( \frac{bh}{m} \right) = \sum_{\psi} c(\psi) \psi(h) \quad \text{for} \quad \gcd(h, m) = 1 \tag{8.9}
\]

with Fourier coefficients

\[
c(\psi) = \frac{1}{\phi(m)} \sum_{j=1}^{m} e\left( \frac{bj}{m} \right) \overline{\psi}(j),
\]

where the prime signalizes that we only consider those \( j \) coprime to \( m \). We note that

\[
c(\psi) = \frac{\psi(b)}{\phi(m)} \sum_{j=1}^{m} e\left( \frac{bj}{m} \right) \overline{\psi}(b) = \frac{\psi(b)G(\overline{\psi})}{\phi(m)}, \tag{8.10}
\]

where

\[
G(\overline{\psi}) = \sum_{j=1}^{m} e\left( \frac{j}{m} \right) \overline{\psi}(j)
\]

is a Gaussian sum. From (8.9) we obtain

\[
\sum_{n=0}^{N-1} e\left( \frac{b}{m} \lambda^n \right) = \sum_{\psi} c(\psi) \sum_{n=0}^{N-1} \psi(\lambda^n)
\]

\[
= N \sum_{\psi(\lambda) = 1} c(\psi) + \sum_{\psi(\lambda) \neq 1} c(\psi) \frac{1 - \psi(\lambda^N)}{1 - \psi(\lambda)}. \tag{8.11}
\]

Setting \( N = \tau \) in (8.11), we get
and then (8.10) implies

\[ \left| \sum_{n=0}^{\tau-1} e\left( \frac{b}{m} \lambda^n \right) \right| \leq \frac{\tau}{\phi(m)} \sum_{\psi(\lambda) = 1} |G(\psi)|. \]

The last sum contains \( \phi(m)/\tau \) terms, and since \( |G(\psi)| < 1 \) for \( \psi \) trivial and \( |G(\psi)| < \sqrt{m} \) for \( \psi \) nontrivial, we obtain

\[ \left| \sum_{n=0}^{\tau-1} e\left( \frac{b}{m} \lambda^n \right) \right| \leq \frac{\tau}{\phi(m)} \left( \frac{\phi(m)}{\tau} - 1 \right) \sqrt{m} = \sqrt{m} - \frac{\tau}{\phi(m)} (\sqrt{m} - 1), \]

and (8.7) is shown. For \( 1 < N < \tau \), we use (8.7), (8.10), (8.11), and (8.12) to get

\[ \left| \sum_{n=0}^{N-1} e\left( \frac{b}{m} \lambda^n \right) \right| \leq N \left| \sum_{n=0}^{\tau-1} e\left( \frac{b}{m} \lambda^n \right) \right| + 2 \sum_{\psi(\lambda) \neq 1} \frac{|c(\psi)|}{|1 - \psi(\lambda)|} \]

\[ \leq N \left( \frac{\sqrt{m}}{\tau} - \frac{\sqrt{m} - 1}{\phi(m)} \right) + \frac{2}{\phi(m)} \sum_{\psi(\lambda) \neq 1} \frac{|G(\psi)|}{|1 - \psi(\lambda)|}. \]

Now \( |G(\psi)| < \sqrt{m} \) and each \( \tau \)th root of unity \( e(j/\tau) \) occurs exactly \( \phi(m)/\tau \) times as a value of \( \psi(\lambda) \), so that

\[ \left| \sum_{n=0}^{N-1} e\left( \frac{b}{m} \lambda^n \right) \right| \leq N \left( \frac{\sqrt{m}}{\tau} - \frac{\sqrt{m} - 1}{\phi(m)} \right) + \frac{2 \sqrt{m}}{\phi(m)} \cdot \frac{\phi(m)}{\tau} \sum_{j=1}^{\tau-1} \left| 1 - e\left( \frac{j}{\tau} \right) \right|^{-1} \]

\[ = N \left( \frac{\sqrt{m}}{\tau} - \frac{\sqrt{m} - 1}{\phi(m)} \right) + \frac{\sqrt{m}}{\tau} \sum_{j=1}^{\tau-1} \csc \frac{\pi j}{\tau}, \]

and (8.8) follows from the argument in [226, p. 574].

In important special cases, the exponential sums in question can actually vanish. A general result in this direction is the following.

**8.4. Theorem (Niederreiter [226]).** Let \( m = p^a \), \( p \) prime, \( a > 1 \), let \( \gcd(b, m) = \gcd(\lambda, m) = 1 \), and let \( z_0, z_1, \ldots \) be a sequence of integers with \( z_{n+1} = \lambda z_n + r \) \((n = 0, 1, \ldots)\) such that the period \( \tau = \tau(m) \) modulo \( m \) satisfies \( \tau = pt(p^{a-1}) \). Then
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\[
\sum_{n=0}^{\tau-1} e\left(\frac{b}{m} z_n\right)e\left(\frac{cn}{\tau}\right) = 0 \text{ for all integers } c \equiv 0 \pmod{p}.
\]

Using this fact, one can then establish the following improvement on (8.3) and (8.5) in the case under consideration.

8.5. **Theorem (Niederreiter [226]).** Suppose the conditions of Theorem 8.4 are satisfied, and let \( \lambda \) belong to the exponent \( \mu \) modulo \( m \). Then,

\[
\left| \sum_{n=0}^{N-1} e\left(\frac{b}{m} z_n\right) \right| < \left(\frac{m\tau}{\mu}\right)^{1/2} \left(\frac{2(p-1)}{\pi p} \log \tau + \frac{3}{4}\right) \text{ for } 1 \leq N \leq \tau.
\]

The estimates in this section are best possible, apart from the logarithmic factors that occur in those cases where a sum over a part of the period is considered. This applies even to the most general estimate (8.3), since it is shown in [224] that there are instances in which this particular sum is at least of the order of magnitude \( m^{1/2} \), while \( \tau = \mu \).

9. **Equidistribution test.** To check the performance of a sequence of LCPRN, we subject it first to test A from §6, the equidistribution test. We will also investigate the behavior of PRN generated by higher-order recurrences with respect to this test. We recall that this test amounts to calculating (or estimating) the discrepancy \( D_N \) of an initial segment \( x_0, x_1, \ldots, x_{N-1} \) of the given sequence of PRN.

If \( x_0, x_1, \ldots, x_{r-1} \) is the full period of a sequence of LCPRN, then the discrepancy \( D_r \) (or \( D_r^* \)) can actually be evaluated in the most interesting cases. For instance, if we have \( r = m \), as happens for \( m = 2^a, \alpha > 2, \lambda \equiv 1 \pmod{4} \), and \( r \) odd (see §7), it is clear that \( D_r = D_r^* = 1/m \). For the homogeneous case with \( m = p^a, p \) an odd prime, \( \alpha > 1, \lambda \) a primitive root modulo \( m \), we get \( D_r^* = 1/m \) (cf. [196]) and \( D_r = 2/m \). Again in the homogeneous case and for \( m = 2^a, \alpha > 3 \), we have \( D_r^* = 3/m \) and \( D_r = 4/m \) if \( \lambda \equiv 5 \pmod{8} \) and \( D_r^* = 5/m \) and \( D_r = 6/m \) if \( \lambda \equiv 3 \pmod{8} \). The first results of this type (cf. [140]) were very weak, with an order of magnitude for \( D_r \) even worse than \( m^{-1/2} \). The fact that the correct order of magnitude is \( m^{-1} \) was pointed out in [218]. See also [57], [58] for special cases.

As was already mentioned in §7, in an actual calculation based on LCPRN we will only use an initial segment of the period since any influence of the periodicity property could prove ruinous. Therefore, the meaningful part of the equidistribution test is the estimation of \( D_N \) for \( N < \tau \). This was first achieved by Niederreiter [221]. Subsequent improvements, simplifications, and generalizations are contained in [226]. All the theorems in this section stem from these two papers. We occasionally incorporate some small improvements.

In the case of a prime modulus, we can use Lemma 3.3 and the estimates in §8 directly, and we note that \( \tau = \mu \) in the inhomogeneous case because of Lemma 7.1.

9.1. **Theorem.** For a sequence of LCPRN with a prime modulus \( m \) we have
for $1 < N < \tau$ in the homogeneous case and

$$D_N < \frac{\sqrt{m}}{N} \left( \frac{2}{\pi} \log \tau + \frac{2}{5} \right) \left( \frac{2}{\pi} \log m + \frac{2}{5} \right)$$

$$+ \left( \frac{\sqrt{m}}{\tau} - \frac{1}{\sqrt{m} + 1} \right) \left( \frac{2}{\pi} \log m + \frac{2}{5} \right) + \frac{1}{m}$$

for $1 < N < \tau$ in the inhomogeneous case.

An estimate based on Corollary 3.2 can also be established. The resulting bounds (cf. [226, Theorems 1 and 4]) are more complicated, but often somewhat better than those in Theorem 9.1, although the order of magnitude remains the same.

For a prime power $m$, say $m = p^\alpha$ with $p$ prime and $\alpha > 2$, there is a difficulty with applying Lemma 3.3 since there will now appear exponential sums for which, in the notation of §8, we have $\gcd(b, m) > 1$. In this case, it is advantageous to use Corollary 3.2 which allows us to "cut out" many of these bothersome sums. We again need the numbers $\kappa$ and $\omega$ introduced prior to Lemma 7.1, as well as one more parameter. With $\mu(q)$ denoting the exponent to which the multiplier $\lambda$ belongs modulo $q$, we define a positive integer $\beta$ as follows: if $p$ is odd, let $\beta$ be the largest integer such that $p^\beta$ divides $\lambda^{\mu(p)} - 1$; if $p = 2$, let $\beta$ be the largest integer such that $2^\beta$ divides $\lambda^{\mu(4)} - 1$. The number $\beta$ is small for the common choices of multipliers. If $p$ is odd and $\lambda$ is a primitive root modulo $m$, then $\beta$ attains its minimal value $\beta = 1$. If $p = 2$, then $\beta$ attains its minimal value $\beta = 2$ for $\lambda \equiv 5 \pmod{8}$, whereas $\beta = 3$ for $\lambda \equiv 3 \pmod{8}$. We set $\gamma = \beta + \omega - \kappa$ and note that we always have $\gamma > 0$. We write $\mu$ for $\mu(m)$.

9.2. Theorem. For a sequence of LCPRN with modulus $m = p^\alpha$, $p$ prime, $\alpha > 2$, which satisfies $\gamma < \alpha$ and

$$p^\gamma < (0.24) \left( \frac{m\tau}{\mu} \right)^{1/2} \left( \frac{2(p - 1)}{\pi p} \log \tau + \frac{3}{4} \right), \quad (9.1)$$

we have

$$D_N < \frac{p^{3/2} - p^{1/2}}{p^{3/2} - 1} X \log \left[ 1 + \frac{4(p^{3/2} - 1)}{(p^{3/2} - p^{1/2})X} \right] + \left( \frac{p^{3/2}}{p^{3/2} - 1} + \frac{\log p}{p} \right) X$$

for $1 < N < \tau$, where

$$X = \frac{4}{\pi N} \left( \frac{m\tau}{\mu} \right)^{1/2} \left( \frac{2(p - 1)}{\pi p} \log \tau + \frac{3}{4} \right).$$

In practical cases, $m$ and $\tau$ are large, so that the condition (9.1) can be
satisfied by choosing the parameters in such a way that $\gamma < \alpha/2$. For the homogeneous case, the above result simplifies somewhat since $\tau = \mu$ and $\omega = \kappa$. Consequently, we have then $\gamma = \beta$, and as the above information about $\beta$ shows, the condition (9.1) is easily satisfied in all reasonable circumstances. In the inhomogeneous case, the common choice $m = 2^\alpha$, $\alpha > 3$, $\lambda \equiv 5 \pmod{8}$, and $r$ odd leads to $\gamma = 3$, and (9.1) holds trivially.

An inspection of the bounds in Theorems 9.1 and 9.2 shows that in terms of orders of magnitude we have $N D_N = O(m^{1/2} \log^2 m)$. Since $N D_N \leq N$ is trivial, the estimates are only of interest for $N$ appreciably larger than $m^{1/2} \log^2 m$, say $N \approx m^{1/2+\varepsilon}$ with $\varepsilon > 0$. But $\tau > N$, and so $\tau$ should be at least of this order of magnitude. If this is so, then the PRN have a small discrepancy and thus pass the equidistribution test. The larger we can choose $N$ (i.e., the larger $\tau$), the better the distribution behavior will be. This provides again support for the familiar rule of thumb that large periods are preferable.

The equidistribution test is not very selective since the results only depend on the values of $\gamma$, $\tau$, and $\mu$. Thus, the test does not allow a distinction between the various primitive roots $\lambda$ we might want to use for an odd prime or prime power modulus in the homogeneous case, or between the various multipliers $\lambda \equiv 5 \pmod{8}$ in the case where $m$ is a power of 2 and $r$ is odd. It is guaranteed, however, that with such choices of parameters the test will be passed with flying colors. We note also that the initial value $y_0$ and the increment $r$ only play an indirect rôle, insofar as they influence the values of $\gamma$ and $\tau$.

The above results can be extended to arbitrary moduli (see [221, §5]). Furthermore, these discrepancy estimates yield effective error bounds for quasi-Monte Carlo integrations based on the nodes $x_0, x_1, \ldots, x_{N-1}$ because of the inequalities in §2. The estimate $N D_N = O(m^{1/2} \log^2 m)$ established above is nearly best possible, as the following lower bounds indicate. We may state these results from [226, §5] in terms of $D_N^\bullet$ since [226, Lemma 11] holds with $D_N$ replaced by $D_N^\bullet$. Because of $D_N^\bullet < D_N$, we automatically get lower bounds for $D_N$ as well.

9.3. Theorem. For any prime modulus $m > 3$, any multiplier $\lambda$ which is a primitive root modulo $m$, and any increment $r$, there exists an initial value $y_0$ with $(\lambda - 1)y_0 + r \equiv 0 \pmod{m}$ such that the derived sequence of LCPRN satisfies $N D_N^\bullet > \frac{1}{8} m^{1/2}$ for some $N$ with $1 < N < \tau$.

The additional condition $(\lambda - 1)y_0 + r \equiv 0 \pmod{m}$, which does not appear in this form in the original result, can be obtained by excluding from the sum $S$ in [226, p. 593] the unique residue $b$ with $(\lambda - 1)b + r \equiv 0 \pmod{m}$ rather than the residue $b = 0$. A similar modification of [226, Theorem 10] yields the following result.

35 Some qualitative evidence for the good performance of LCPRN under the equidistribution test was already presented by Franklin [71] who set up a continuous model and showed that for any integer $\lambda > 1$ and any real $\theta$ the sequence defined by $x_{n+1} = (\lambda x_n + \theta)$, $n = 0, 1, \ldots$, is uniformly distributed in $I$ for almost all initial values $x_0 \in I$ (in the sense of Lebesgue measure).

See also [248, Chapter 3] for a detailed discussion of the case $\theta = 0$.
9.4. Theorem. Let \( m = p^a \), \( p \) prime, \( a \geq 2 \), be a modulus, let the multiplier \( \lambda \) belong to the largest possible exponent modulo \( m \), and let \( r \) be an increment. If \( p \) is odd, there exists an initial value \( y_0 \) with \( (\lambda - 1)y_0 + r \not\equiv 0 \pmod{p} \) such that the derived sequence of LCPRN satisfies

\[
ND_N^* \geq \frac{(p^2 - 1)^{1/2}}{8p} \cdot m^{1/2}
\]

for some \( N \) with \( 1 \leq N \leq \tau \). If \( p = 2 \) and \( r = 0 \), there exists an odd initial value \( y_0 \) such that the derived sequence of LCPRN satisfies

\[
ND_N^* \geq \frac{1}{8\sqrt{2}} \cdot m^{1/2}
\]

for some \( N \) with \( 1 \leq N \leq \tau \). If \( p = 2 \) and \( r \) odd, there exists an initial value \( y_0 \) such that the derived sequence of LCPRN satisfies \( ND_N^* \geq \frac{1}{8} \cdot m^{1/2} \) for some \( N \) with \( 1 \leq N \leq \tau \).

It is remarkable how close the discrepancy for parts of the period of a sequence of LCPRN is to that of a randomly chosen sequence (compare with §3).

The techniques employed in this section can also be applied to sequences of PRN generated by higher-order recurrences. Using the notation and the result of Theorem 8.1 as well as Lemma 3.3, we obtain the following information concerning the equidistribution test.

9.5. Theorem. For any prime \( p \) and any sequence of PRN derived from a \( k \)th order linear recurrence relation (8.1) with \( a_0 \not\equiv 0 \pmod{p} \), we have

\[
D_N < \frac{1}{p} + \frac{1}{N} \left( \frac{p^k}{\mu} \right)^{1/2} \left( \frac{2}{\pi} \log \tau + \frac{2}{5} \right) \left( \frac{2}{\pi} \log p + \frac{2}{5} \right)
\]

\[
+ \frac{1}{\tau} \left( \frac{p^k \tau - \tau^2}{\mu} \right)^{1/2} \left( \frac{2}{\pi} \log p + \frac{2}{5} \right)
\]

for \( 1 \leq N \leq \tau \). \hspace{1cm} (9.2)

9.6. Corollary. For any \( k \)th order Tausworthe generator modulo the prime \( p \) we have

\[
D_N < \frac{1}{p} + \frac{p^{k/2}}{N} \left( \frac{2}{\pi} \log \tau + \frac{2}{5} \right) \left( \frac{2}{\pi} \log p + \frac{2}{5} \right) + \frac{1}{\tau} \left( \frac{2}{\pi} \log p + \frac{2}{5} \right)
\]

for \( 1 \leq N \leq \tau = p^k - 1 \). \hspace{1cm} (9.3)

The bound (9.2) is only meaningful if \( N \), and so the period \( \tau \), is considerably larger than \( p^{k/2} \). Thus, as in the first-order case, recurrences producing long periods are preferable. If \( N \geq p^{(k/2)+1+\varepsilon} \), \( \varepsilon > 0 \), then the main term in (9.3) is \( 1/p \). This term is needed since all PRN are rationals with denominator \( p \), so that we must have \( D_N \geq 1/p \). On the other hand, the upper bound (9.3) differs then by very little from this trivial lower bound implied by the discreteness of the sequence, which confirms the notion that Tausworthe generators, say of order \( k \geq 3 \), lead to an extremely even distribution. For small values of \( k \) the second term on the right-hand side of
(9.3) is required, at least up to the logarithmic factors, because of the following result: for every primitive polynomial modulo $p$ of degree $k$ there exists a corresponding Tausworthe generator such that $ND_N^* > \frac{1}{8} p^{k/2}$ for some $N$ with $1 < N < \tau$ (see [226, p. 596]).

10. Interdependence of successive terms. The results of the preceding section have shown that a sequence of LCPRN will pass the equidistribution test as soon as its period is large. To achieve a finer distinction between the various multipliers, more powerful tests have to be applied. In particular, one has to take a closer look at order properties and relations between successive terms in the sequence.

Although numerical data about the performance of various LCPRN under stringent tests had already been collected earlier (cf. [143], [332], and the survey article [135]), the first effective theoretical results were only obtained in the 1960s when the problem of calculating serial correlation coefficients (see §6, test E) was solved satisfactorily. Coveyou [46] considered the question in the context of a continuous model, whereas Greenberger [86] provided a good estimate for serial correlation coefficients. To minimize the resulting expressions, the choice $\lambda \approx m^{1/2}$ suggested itself, leading to proposals such as $m = 2^{35}, \lambda = 2^{18} + 3$. A few years later, Jansson [141] gave exact formulas for serial correlation coefficients. All these results refer to the full period and to cases in which either all residues or a known set of residues modulo $m$ are generated. See also [1], [56], [60], [142, Chapter 6], [154, §3.3.3], [327]. The values of serial correlation coefficients can be conveniently expressed in terms of generalized Dedekind sums. An allusion to these sums occurs already in [86]. Algorithms for the calculation of serial correlation coefficients based on the reciprocity law for generalized Dedekind sums are presented in [56], [60], [154, §3.3.3]. On the basis of such calculations and an analogy with diophantine approximations, it is suggested in [1] that for the homogeneous case and $m$ a power of 2, one should select a multiplier $\lambda$ with $\lambda \equiv 5 \pmod{8}$ and $
.lambda \approx m\xi/4$, where $\xi = (\sqrt{5} - 1)/2$ is the golden section number, in order to get small serial correlation.

Numerical data on the run test (see §6, test C) for commonly used generators were already collected in the 1950s (cf. [332]). This test has also been applied to a special class of LCPRN obtained by taking the Mersenne prime $m = 2^{31} - 1$ as the modulus in the multiplicative congruential method, with the multiplier $\lambda$ being a primitive root modulo $m$. This proposal is due to D. H. Lehmer and has the merit of allowing a quite convenient computer implementation (cf. [91], [183]). Particularly good results were achieved with $\lambda = 7^s$ (cf. [181]), but there can be problems with multipliers that are too small (cf. [67]). For further work on these generators, see [60], [61], [136], [239], [286].

A feasible implementation of the spectral test (see §6, test F) offers some challenges, so that numerical data were only recently compiled in a systematic form (for some sample calculations, see [154, §3.3.4]). We refer to [55] for dimensions $2 < s < 4$ and [107] for dimensions $2 < s < 5$.

Because of the simple nature of the recursion (7.1), it is clear that sequences of LCPRN will be endowed with some intrinsic structure. This was made
explicit by the results of Marsaglia [186], [187] on the lattice (or "crystalline")
structure induced by \( s \)-tuples of successive LCPRN. For notational reasons, it
is more convenient here to work with the sequence of integers \( y_0, y_1, \ldots \)
generated by (7.1). We choose a dimension \( s \geq 2 \) and consider the lattice
points \( y_n = (y_n, y_{n+1}, \ldots, y_{n+s-1}) \), \( n = 0, 1, \ldots, \tau - 1 \), where \( \tau \) is, as usual,
the period of the sequence. Then we subtract from each of those points the
vector \( \mathbf{c} = (c_0, c_1, \ldots, c_{s-1}) \) whose coordinates result from (7.1) by starting
with \( c_0 = 0 \). In this way we obtain \( \tau \) lattice points \( y_0, y_1, \ldots, y_{\tau-1} \). We now
inspect the crystalline structure of these points; it may happen, for instance,
that they all lie on a rather "coarse" lattice.\(^{36}\) A measure for the "coarseness"
of a lattice is its unit-cell volume, obtained by calculating the absolute value
of a determinant whose row vectors form a basis for the lattice. Marsaglia
[187] shows that for any choice of parameters in (7.1) the points \( y_0, y_1, \ldots, y_{\tau-1} \) all lie on a lattice with unit-cell volume \( m^{s-1} \), which he finds
unacceptably large. However, the result is stated this way for greater effect.
We must not forget that the PRN themselves are produced by dividing each
\( y_n \) by \( m \). Taking into account the resulting scaling factor \( m^{-s} \), the actual
unit-cell volume becomes \( m^{-1} \), which is about \( 2^{-35} \) for the standard genera-
tors and thus quite reasonable.

Marsaglia [188, p. 270] promises methods for perturbing congruential
generators which would prevent the above points from lying on lattices with
unit-cell volume greater than 1. This can of course be achieved in a trivial
manner. For instance, in the homogeneous case we take the sequence \( y_0, y_1, \ldots \)
with initial value \( y_0 = 1 \) and throw in the numbers \( y_{-s+1} = y_{-s+2} = \ldots = y_{-1} = 0 \), ignoring the recursion for these indices. Then the vectors
\( y_{-s+1}, y_{-s+2}, \ldots, y_0 \) span a parallelepiped of volume 1 and the desired
property holds. But such juggling acts are evidently not worth the effort since
it is hard to see why the new sequence should be closer to randomness than
the original one.

This raises some fundamental doubts about whether the "lattice test" (to
use Marsaglia's term) has the proper credentials of a meaningful statistical
test. As the above example indicates, it does not satisfy a mandatory
requirement on such tests, namely that of stability. A small perturbation of
the data (such as changing some terms slightly or adding a small number of
new values) should not affect the outcome in any significant way, and this is
clearly the case in all the tests mentioned in §6A-G. The "lattice test", on the
other hand, reflects purely arithmetic coincidences that can be changed on a
whim.

Beyer [19] stressed the fact that the configuration of points \( y_0, y_1, \ldots, y_{\tau-1} \), when extended with period \( m \) in each coordinate to all of \( \mathbb{R}^s \),
need not form a lattice by itself. This will be true, however, if \( \tau = m \), which
happens in the inhomogeneous case for a suitable choice of parameters (see
§7). In the homogeneous case with a multiplier belonging to the largest
possible exponent modulo \( m \), the periodically extended configuration can be
viewed as the union of a limited number of translated versions of a fixed

\(^{36}\) A lattice in \( \mathbb{R}^s \) is a set of lattice points obtained by forming all integral linear combinations
of \( s \) linearly independent lattice points.
lattice. This has led to attempts at analyzing this lattice with the aim of finding criteria for "optimal" generators. The viewpoint commonly espoused in this connection is that a generator is the better the closer its associated lattice is to a "cubic" lattice (cf. [20], [188]). This rather ill-advised notion is based on simplistic myths about the cubic lattice. In the realm of facts, the cubic lattice is of dubious merit. Consider the \( N = m^s \) points \((n_1/m, \ldots, n_s/m)\), with the \( n_i \) running independently through the integers 0, 1, \ldots, \( m - 1 \). They form a pleasing arrangement on a cubic grid. However, since no point falls into the interval \((0, 1/m) \times I^s \) of volume \( m^{-1} \), the discrepancy \( D_N \) of these points satisfies \( D_N > m^{-1} = N^{-1/s} \), a pitiable performance. Even more striking evidence for the fallacy of the "cubic-lattice criterion" will be presented in \( \S 11 \), where we will prove on the basis of effective theoretical results that a multiplier which Marsaglia [188, p. 275] calls "a candidate for the best of all multipliers" because of the closeness of its 2-dimensional lattice to a cubic lattice actually shows a miserable behavior under the 2-dimensional serial test.

There is a way in which the lattice structure can be profitably studied to gain insight into the behavior of LCPRN. It is based on the observation of Marsaglia [186] that the points \( y_0, y_1, \ldots, y_{t-1} \) all lie on a limited number of parallel hyperplanes. If this number is too small, then there are large portions of the unit cube \( I^t \) devoid of such points, and thus the generator is unacceptable. Equivalently, one may measure the distances between neighboring hyperplanes containing these points, and if the distances are too large, the generator should be disqualified. In the latter form, this test is very strongly linked with the spectral test (cf. [154, p. 100]). In fact, the "hyperplane test" may be viewed as a geometric version of the spectral test.

Relations between successive terms have also been investigated in the case of Tausworthe generators (see \( \S 7 \)). Results about the serial correlation can already be found in the early papers on these generators ([333], [362]). The run test was applied in [337] and the weaker version of the serial test (see \( \S 6 \), test G, second part) in [68], [333]. In [338] excellent multidimensional uniformity properties are reported for a 3-term recurrence of order 607 modulo 2. But, as pointed out in [363], Tausworthe generators also have to be used with care and the desired degree of statistical independence should be ascertained before adopting them. The lattice structure of Tausworthe generators was studied in [89].

11. Serial test. The discussion in \( \S 10 \) has not included the strongest test for the statistical independence of successive terms, which will now be analyzed in detail. To set up the serial test (see \( \S 6 \), test G) for a sequence \( x_0, x_1, \ldots \) of LCPRN, we choose a dimension \( s > 2 \) and consider the points \( x_n = (x_n, x_{n+1}, \ldots, x_{n+s-1}) \), \( n = 0, 1, \ldots, \) in \( I^s \) and their distribution behavior, measured quantitatively by the discrepancy \( D_{IN} \). In the easier case of the multiplicative congruential method, the points \( x_n \) may be represented explicitly by the compact formula

\[
x_n = \{ \lambda^n x_0 \lambda \} \quad \text{for } n = 0, 1, \ldots ,
\]

where \( \lambda \) denotes the lattice point \((1, \lambda, \lambda^2, \ldots, \lambda^{s-1}) \in \mathbb{Z}^s \). It is therefore not surprising that this lattice point should play a special rôle in our
investigation, and it will also transpire that the very same lattice point is of equal importance in the mixed congruential method.

The first attempt at getting a grip on the multidimensional distribution properties of LCPRN was undertaken in a paper of Franklin [72]. Here a continuous model was developed by accepting as an initial value any number $x_0$ in the unit interval $I$ and generating a sequence $x_0, x_1, \ldots$ by the recursion $x_{n+1} = \{\lambda x_n + \theta\}$, $n = 0, 1, \ldots$, where the integer $\lambda > 1$ and the real number $\theta$ are fixed. For a given dimension $s \geq 2$, we associate with this sequence the points $x_0, x_1, \ldots$ in $I^s$ in the same way as above. It is easily seen that the sequence $x_0, x_1, \ldots$ can never be uniformly distributed in $I^s$. However, the sequence is, in a certain sense, almost always “asymptotically” uniformly distributed. To make this explicit, we have to adopt a more careful notation. We fix $\theta$ and emphasize the dependence on $\lambda$ by writing $x_n(\lambda)$ instead of $x_n$. Since the initial value $x_0$ is independent of $\lambda$, we may retain the simpler symbol for it. Franklin shows now that for almost all $x_0 \in I$ (in the sense of Lebesgue measure) we have

$$\lim_{\lambda \to \infty} \lim_{N \to \infty} \frac{1}{N} \sum_{n=0}^{N-1} c_J(x_n(\lambda)) = |J|$$

for all subintervals $J$ of $I^s$, where $c_J$ is the characteristic function and $|J|$ the Lebesgue measure of $J$. Furthermore, if $f$ is a continuous function on $I^s$, then for almost all $x_0 \in I$ the relation

$$\lim_{\lambda \to \infty} \lim_{N \to \infty} \frac{1}{N} \sum_{n=0}^{N-1} f(x_n(\lambda)) = \int_{I^s} f(t) \, dt$$

is satisfied (see also [304] for related results). The regularity condition on $f$ may be relaxed. For the case $\theta = 0$ corresponding to the multiplicative congruential method, Ermakov [65] established a quantitative refinement in the form of an “asymptotic” central limit theorem. Let $f$ again be a continuous function on $I^s$ and denote by

$$R_N(f, x_0, \lambda) = \frac{1}{N} \sum_{n=0}^{N-1} f(x_n(\lambda)) - \int_{I^s} f(t) \, dt$$

the integration error. Then we have

$$\lim_{N \to \infty} \lim_{\lambda \to \infty} \left| \left\{ x_0 \in I : R_N(f, x_0, \lambda) < \frac{\sigma u}{\sqrt{N}} \right\} \right| = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{u} e^{-t^2/2} \, dt,$$

(11.1)

where the constant $\sigma$ designates a certain standard deviation depending on $f$. The condition on $f$ may be relaxed in the same way as in Franklin’s theorem. These results suggest that, at least on probabilistic grounds, things should work out all right for a sequence of LCPRN with large modulus and large multiplier, but they do not provide any information about individual sequences of LCPRN. We note that the shuffling scheme of MacLaren and Marsaglia [184] was analyzed by Rosenblatt [260] for dimensions $s = 2$ and $3$ in terms of an analogous probabilistic model.
It is, of course, considerably more difficult to get effective results about specific sequences of LCPRN. Recent research of the author has produced a satisfactory theory for any dimension $s > 2$. We let $x_0, x_1, \ldots$ be an arbitrary sequence of LCPRN and define the points $x_0, x_1, \ldots$ in $I^s$ as before. To recognize clearly the dependence on the dimension, we denote by $D_N^{(s)}$ the discrepancy of the points $x_0, x_1, \ldots, x_{N-1}$ in $I^s$. It will be convenient to use an expression that is closely related to the one introduced in Definition 4.3.

11.1. Definition. For a modulus $m$, a multiplier $\lambda$, a dimension $s > 2$, and a positive integer $d$ we set

$$R^{(s)}(\lambda, m, d) = \sum_{\substack{h \pmod{m} \atop h \cdot \lambda \equiv 0 \pmod{d}}} r(h)^{-1},$$

where $\lambda = (1, \lambda, \lambda^2, \ldots, \lambda^{s-1}) \in \mathbb{Z}^s$, the summation symbol is the same as in (4.5), and $r(h)$ is defined by (4.3).

If $d = m$, then $R^{(s)}(\lambda, m, m)$ is the same as $R(\lambda, m)$ from Definition 4.3. The subsequent estimates for $D_N^{(s)}$ will be in terms of $R^{(s)}(\lambda, m, d)$, where $d$ depends in a known way on the parameters determining the specific sequence. The case $N = r$, i.e., when we consider the full period, is somewhat easier and will be treated first. The method rests on Lemma 3.9 and the estimates in §8. We simplify the upper bounds by using the inequality (4.19). For a prime modulus $m$ we obtain the following.

11.2. Theorem (Niederreiter [227], [229]). For a sequence of LCPRN with prime modulus $m$ we have

$$D_r^{(s)} \leq \frac{s}{m} + \frac{1}{\tau} (m - \tau)^{1/2} \left( \frac{2}{\pi} \log m + \frac{7}{5} \right)^s + \frac{1}{2} R^{(s)}(\lambda, m, m).$$

The second term in this upper bound is nonincreasing as a function of $\tau$ and so becomes minimal for the maximal value of $\tau$, namely $\tau = m - 1$. Therefore, if we choose $\lambda$ to be a primitive root modulo $m$, then we obtain

$$D_{m-1}^{(s)} \leq \frac{s}{m} + \frac{1}{m - 1} \left( \frac{2}{\pi} \log m + \frac{7}{5} \right)^s + \frac{1}{2} R^{(s)}(\lambda, m, m). \quad (11.2)$$

To get LCPRN for which $s$ successive terms have a low degree of statistical dependence, we now select a primitive root $\lambda$ modulo $m$ yielding a small value of $R^{(s)}(\lambda, m, m)$. Since $R^{(s)}(\lambda, m, m) = R(\lambda, m)$, this is equivalent to the desideratum that $\lambda$ be a good lattice point modulo $m$ (compare with §4). We have thus established an intriguing connection between pseudo-random number generation and the theory of good lattice points. As it stands, this theory does not indicate how small we can make $R^{(s)}(\lambda, m, m)$ since the requirement that $\lambda$ be a primitive root modulo $m$ is of no consequence there. It is therefore necessary to prove an existence theorem going beyond those in §4 (cf. [229]). In combination with (11.2), one obtains then the following result.

11.3. Theorem (Niederreiter [229]). For any prime modulus $m$ and any dimension $s > 2$, there exists a primitive root $\lambda$ modulo $m$ such that
\begin{equation}
D_{m-1}^{(s)} < \frac{1}{m-1} \left(1 + \frac{(m-2)(s-1)}{\phi(m-1)} \right) \left( \frac{2}{\pi} \log m + \frac{7}{5} \right), \tag{11.3}
\end{equation}

where \( \phi \) is Euler's totient function.

Since it is well known that \( \phi(b)^{-1} = O(b^{-1} \log \log b) \) with an effective implied constant, we conclude that for a prime modulus \( m \) the parameters in (7.1) can be chosen in such a way that \( D_{m-1}^{(s)} = O(m^{-1}(\log m)^s \log \log m) \) with a constant only depending on \( s \). This should be compared with the result from §3 according to which the best distribution of \( m-1 \) points in \( I^s \) currently known involves a discrepancy of the order of magnitude \( m^{-1}(\log m)^{s-1} \). Therefore, sequences of appropriately selected LCPRN perform remarkably well under the serial test. Actually, the bound in (11.3) represents the expected magnitude of \( D_{m-1}^{(s)} \) for a random primitive root modulo \( m \). This conforms with an experience encountered by people working with empirical data, namely that a blindly chosen multiplier will tend to lead to good statistical properties.

In the case of a prime power modulus, the bounds for \( D_{r}^{(s)} \) become even simpler, although we have to employ now the general expression introduced in Definition 11.1. We also need the number \( y \) defined prior to Theorem 9.2.

11.4. Theorem (Niederreiter [230], [233]). For a sequence of LCPRN with modulus \( m = p^\alpha \), \( p \) prime, \( \alpha > 2 \), and satisfying \( \gamma < \alpha \) we have

\begin{equation}
D_{r}^{(s)} < \frac{s}{m} + \frac{1}{2} R^{(s)}(\lambda, m, p^{a-y}). \tag{11.4}
\end{equation}

For the homogeneous case this was already shown in [227], [229]; then we have of course \( \gamma = \beta \) (see §9). It is perhaps worthwhile to list the values of \( \gamma \) for \( m = 2^\alpha, \alpha > 3 \), and the most common choices of the multiplier \( \lambda \) and the increment \( r \). In the homogeneous case, \( \lambda \equiv 5 \pmod{8} \) leads to \( \gamma = 2 \) and \( \lambda \equiv 3 \pmod{8} \) has \( \gamma = 3 \). In the inhomogeneous case with \( r \) odd, \( \lambda \equiv 5 \pmod{8} \) yields \( \gamma = 0 \), whereas \( \lambda \equiv 3 \pmod{8} \) leads to \( \gamma = 2 \). Since small values of \( \gamma \) are preferable, this provides another reason why \( \lambda \equiv 5 \pmod{8} \) and \( r \) odd seems to be a good choice. On the whole, however, the controversy about whether the multiplicative or the mixed congruential method is superior is rather academic since the difference between the two cases is negligible if Theorem 11.4 is taken as indicative. There is other evidence for the essential equivalence of the two methods, and it can be safely said that a consensus has been reached on this issue (compare with [57, p. 857], [154, p. 20], [188, p. 251]).

Because of the connection with good lattice points, and also for aesthetic reasons, a "symmetric" expression of the form \( R^{(s)}(\lambda, m', m') \) is more agreeable than the one appearing in (11.4). Fortunately enough, the latter expression is related in a fairly simple manner to one of the symmetric type. In fact, for \( 0 < \gamma < \alpha \) we have

\begin{equation}
R^{(s)}(\lambda, m, p^{a-y}) < (1 + 2\gamma \log p)^s R^{(s)}(\lambda, p^{a-y}, p^{a-y}) + (1 + (2\gamma p^y \log p)/m)^s - 1 \tag{11.5}
\end{equation}

for \( p \) odd, according to [229, Equation (5.12)], and an almost identical result.
holds for \( p = 2 \) (cf. [229, Equation (5.20)]). In the cases of interest, \( \gamma \) is very small, and then \( R^{(\gamma)}(\lambda, m, p^{a-\gamma}) \) is of the same order of magnitude as \( R^{(\gamma)}(\lambda, p^{a-\gamma}, p^{a-\gamma}) \). This amounts to saying that the arithmetic properties of \( \lambda \) are not so much relevant modulo \( m = p^a \) but modulo \( p^{a-\gamma} \).

We turn now to estimates for the discrepancy \( D_N^{(\gamma)} \) referring to parts of the period. The following result can be shown for a prime modulus.

11.5. **Theorem (Niederreiter [230], [233]).** For a sequence of LCPRN with prime modulus \( m \) we have

\[
D_N^{(\gamma)} < \frac{s}{m} + \frac{\sqrt{m}}{N} \left( \frac{2}{\pi} \log \tau + \frac{7}{5} \right) \left( \frac{2}{\pi} \log m + \frac{7}{5} \right) + \frac{1}{2} R^{(\gamma)}(\lambda, m, m)
\]

for \( 1 < N < \tau \). (11.6)

If one chooses a primitive root \( \lambda \) modulo \( m \) according to Theorem 11.3, then the third term on the right-hand side of (11.6) is \( O(m^{-1}(\log m)^{\gamma}) \log \log m \), and so \( ND_N^{(\gamma)} = O(m^{1/2}(\log m)^{\gamma+1}) \). We observe that in general we should not expect a smaller exponent of \( m \) for large \( N \) because of Ermakov's central limit theorem in (11.1). If \( N \) is comparable with \( m \) in size, then the order of magnitude of \( ND_N^{(\gamma)} \) resembles that of a "random" sequence of points in \( I^s \) (see \$3).

For the statement of the estimate in the case of a prime power modulus we use the same notation as in Theorem 9.2.

11.6. **Theorem (Niederreiter [230], [233]).** For a sequence of LCPRN with modulus \( m = p^a, p \) prime, \( \alpha > 2 \), and satisfying \( \gamma < \alpha \) we have

\[
D_N^{(\gamma)} < \frac{s}{m} + \frac{1}{N} \left( \frac{m \tau}{\mu} \right)^{1/2} \left( \frac{2(p-1)}{\pi p} \log \tau + \frac{3}{4} \right) \left( \frac{2}{\pi} \log m + \frac{7}{5} \right) + \frac{1}{2} R^{(\gamma)}(\lambda, m, p^{a-\gamma}) \quad \text{for } 1 < N < \tau.
\]

In the homogeneous case, (11.7) can be simplified somewhat on the basis of the identity \( \tau = \mu \). For the mixed congruential method with \( m = 2^a, \alpha > 3 \), and \( r \) odd, we have \( \tau/\mu = 4 \) for \( \lambda \equiv 5 \) (mod 8) and \( \tau/\mu = 2 \) for \( \lambda \equiv 3 \) (mod 8). By using (11.5), we may, up to a constant, replace the last term in (11.7) by \( \bar{R}^{(\gamma)}(\lambda, p^{a-\gamma}, p^{a-\gamma}) \) in the cases of interest.

As we already observed in the discussion of the theory of good lattice points, a quantity such as \( R^{(\gamma)}(\lambda, m, m) \) is awkward to deal with numerically. Consequently, we adopt the same remedy as there and introduce an integer \( \rho^{(\gamma)}(\lambda, m) \) by setting

\[
\rho^{(\gamma)}(\lambda, m) = \rho(\lambda, m).
\]

where \( \lambda = (1, \lambda, \lambda^2, \ldots, \lambda^{s-1}) \in \mathbb{Z}^s \) and \( \rho(\lambda, m) \) is defined by (4.8). We can then rephrase Theorem 4.5 as follows.

11.7. **Lemma.** For any dimension \( s > 2 \) and any integers \( m > 2 \) and \( \lambda \) we have

\[
R^{(\gamma)}(\lambda, m, m) < C_s (\log m)^s / \rho^{(\gamma)}(\lambda, m)
\]
with a constant $C_s$ only depending on $s$.

It follows from the above considerations that a sequence of LCPRN passes the $s$-dimensional serial test if, in the case of a prime modulus $m$, we choose a multiplier $\lambda$ which is a primitive root modulo $m$ and yields a large value of $\rho^{(s)}(\lambda, m)$, and in the case of a prime power modulus $m = p^\alpha$ we select the parameters in such a way that both the period and $\rho^{(s)}(\lambda, p^\alpha)$ are large. That the size of the number from (11.8) is really the correct indicator for the performance of a sequence of LCPRN under the $s$-dimensional serial test is shown by the following result.

11.8. Theorem (Niederreiter [230], [233]). For any sequence of LCPRN with modulus $m$ and multiplier $\lambda$, we have

$$D^{(s)}_N > C'_s / \rho^{(s)}(\lambda, m)$$

for $1 < N < \tau$, where $C'_s$ is a positive constant only depending on the dimension $s$.

We may take $C'_s = s^{-s}$ for $2 < s < 3$ and $C'_s = (\pi/2)(\pi + 1/2)^{-s}$ for $s > 4$. The lower bound may be improved in some important special cases. For instance, if we consider the multiplicative congruential method with $m = 2^\alpha$, $\alpha > 3$, and $\lambda \equiv 5 \pmod{8}$, then

$$D^{(s)}_N > C'_s / \rho^{(s)}(\lambda, 2^{\alpha-2})$$

for $1 < N < \tau$.

For the most popular generators and the serial test for the full period, we can summarize all this information in a more succinct form by employing Vinogradov’s notation $\ll$ to assimilate constants only depending on $s$. If we use a prime modulus $m$ and a primitive root $\lambda$ modulo $m$ in either a multiplicative or a mixed congruential method, or if we use a mixed congruential method with $m = 2^\alpha$, $\alpha > 3$, and $\lambda \equiv 5 \pmod{8}$, and $\tau$ odd, then we obtain

$$\frac{1}{\rho^{(s)}(\lambda, m)} \ll D^{(s)}_\tau \ll \frac{(\log m)^s}{\rho^{(s)}(\lambda, m)}.$$

(11.9)

If we use a multiplicative congruential method with $m = 2^\alpha$, $\alpha > 3$, and $\lambda \equiv 5 \pmod{8}$, then we have

$$\frac{1}{\rho^{(s)}(\lambda, 2^{\alpha-2})} \ll D^{(s)}_\tau \ll \frac{(\log m)^s}{\rho^{(s)}(\lambda, 2^{\alpha-2})}.$$

(11.10)

Thus the order of magnitude of $D^{(s)}_\tau$ is, up to logarithmic factors, completely determined by the order of magnitude of $\rho^{(s)}(\lambda, m)$ in (11.9) resp. $\rho^{(s)}(\lambda, 2^{\alpha-2})$ in (11.10). We take this as the justification for referring, in general, to the integer $\rho^{(s)}(\lambda, m)$ defined by (11.8) as the $(s$-dimensional) figure of merit of the multiplier $\lambda$ (with respect to $m$).

We have with deliberate care emphasized in the notation that the figure of merit $\rho^{(s)}(\lambda, m)$ depends strongly on the dimension $s$. It is implicit in this fact that a multiplier which is excellent for a certain dimension may be unaccept-
able for another dimension. The reason behind this phenomenon has been discussed in §4, and we shall present concrete examples later on. There is thus a relativity principle in operation, according to which the choice of an excellent (or even optimal) multiplier has to be made relative to the dimensionality of, or the desired number of statistically independent successors in, the Monte Carlo problem at hand. There is no such thing as a universally optimal multiplier. Every individual pseudo-random number generator becomes unsuitable for sufficiently high dimensions, even though it may perform well for small dimensions. This can actually be put in a concrete form. For suppose we generate a sequence of LCPRN with modulus \( m \) and multiplier \( \lambda \) and apply the \( s \)-dimensional serial test with \( s > \log_2(m + 1) \). Consider all lattice points \( h = (h_1, \ldots, h_s) \) with the \( h_j \) taking independently the values 0 and 1. There are \( 2^s \) such lattice points, and since \( 2^s > m + 1 \), there exist two distinct lattice points \( h^{(1)}, h^{(2)} \) among these such that \( h^{(1)} \cdot \lambda \equiv h^{(2)} \cdot \lambda \pmod{m} \). Then \( (h^{(1)} - h^{(2)}) \cdot \lambda \equiv 0 \pmod{m} \) and \( r(h^{(1)} - h^{(2)}) = 1 \), and so \( \rho^{(s)}(\lambda, m) = 1 \) for any multiplier \( \lambda \), which is unacceptable. For instance, the modulus \( m = 2^{35} \) certainly has to be ruled out for dimensions \( s > 36 \). In practice, the threshold is of course considerably lower, and this particular modulus is probably dubious for \( s > 10 \). However, this should not be construed as a suggestion that Lehmer’s method breaks down for moderately high dimensions. On the contrary, results such as Theorem 11.3 guarantee that if we are willing to use a sufficiently large modulus, then there is always a sequence of LCPRN which performs well under the \( s \)-dimensional serial test for an arbitrarily given dimension \( s \). Deficiencies that may occur are therefore not the fault of the generator, but of the limited computer capacity. These limits are of course constantly pushed upwards.

As to the variability of excellent multipliers with changing dimension, something of a positive nature can be said. It is based on the inequality (4.12) which we can rewrite as

\[
\rho^{(t)}(\lambda, m) > \rho^{(s)}(\lambda, m) \quad \text{for dimensions } 2 < t < s. \quad (11.11)
\]

We infer from (11.11) that if a multiplier \( \lambda \) is favorable for the dimension \( s \), i.e., if it leads to a large value of \( \rho^{(s)}(\lambda, m) \), then we can expect this multiplier to show an acceptable (though not necessarily optimal) behavior for all lower dimensions. Thus, if a generator is to be used for several purposes, it suffices to choose it in such a way that it satisfies the most stringent statistical independence condition desired in these applications. Viewed from a different angle, the inequality (11.11) tells us that if a sequence of LCPRN fails the serial test for a certain dimension, it will fail the test for all higher dimensions, which is intuitively quite obvious.

We observed in §4 that for the case \( s = 2 \) more information is available in the theory of good lattice points because of intimate connections with the continued fraction algorithm. The same holds for the 2-dimensional serial test, i.e., for the distribution of pairs of successive LCPRN. In the first place, the 2-dimensional figure of merit \( \rho^{(2)}(\lambda, m) \) is related in a very simple manner with a number arising from the continued fraction expansion of the rational \( \lambda/m \). This is obtained from (4.14) which can be rewritten in the form
\[ m / (K + 2) < \rho^{(2)}(\lambda, m) < m / K, \]  
(11.12)

where \( K = K(\lambda / m) \) is defined by (4.13) and represents the maximal partial quotient in the continued fraction expansion of \( \lambda / m \). Therefore, the figure of merit \( \rho^{(2)}(\lambda, m) \) can only be large if \( \lambda / m \) yields small partial quotients. This raises the question of how small we can make \( K(\lambda / m) \) given the modulus \( m \). The number \( K_m \) from §4 pertains to this question, but it is not necessarily appropriate here since we also have to take into account the fact that \( \lambda \) should produce a sequence of LCPRN with large period. For instance, in the case of a prime modulus \( m \) we actually have to consider

\[ P_m = \min_{\lambda} K\left( \frac{\lambda}{m} \right), \]

where \( \lambda \) runs through all primitive roots modulo \( m \). A theoretical estimate for \( P_m \) is available since it was shown in [229, §4] that for any prime \( m \) there exists a primitive root \( \lambda \) modulo \( m \) with

\[ \rho^{(2)}(\lambda, m) \gg \frac{\phi(m - 1)}{\log \phi(m - 1)}, \]

which together with (11.12) implies that

\[ P_m \ll \frac{m \log \phi(m - 1)}{\phi(m - 1)} \ll (\log m) \log \log m. \]

This bound is certainly not best possible.

The 2-dimensional case also permits us to establish finer estimates for the discrepancy \( D^{(2)}_r \) of the full period, at least for the generators of practical interest. This is achieved by a direct approach and was first carried out by Dieter [57], [58]. The method of Dieter is based on an explicit representation of the local deviations \( A(J; \tau) - \tau|J| \) for subintervals \( J \) of \( I^2 \) in terms of generalized Dedekind sums. These sums can then be calculated by means of a reciprocity law and other arithmetic properties, and the resulting algorithm resembles the continued fraction algorithm (or, equivalently, the Euclidean algorithm). In this way, one obtains exact formulas for the various local deviations. We get information about \( D^{(2)}_r \) by establishing a global bound on the local deviations. This can be done using either standard continued fractions (cf. [155]) or continued fractions to nearest integers, the alternative chosen by Dieter. Generalized Dedekind sums can also be applied to get a handle on the permutation test (see §6, test D) for triples ([58], [59]).

Apart from the fact that it only works for the full period, there is a further restriction on the above technique insofar as we have to know exactly which residues modulo \( m \) are generated by the recursion (7.1). This explains the additional conditions in the subsequent result. The bound for \( D^{(2)}_r \) will be in terms of partial quotients arising from a continued fraction algorithm to nearest integers.\(^{38}\) The new feature here is that these partial quotients may also be negative.

11.9. **THEOREM (DIETER [57], [58]).** For a sequence of LCPRN with modulus

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\(^{38}\) See [241, Chapter 5] for the theory of such continued fraction algorithms.
multiplier $\lambda$, and either (i) $\tau = m$; or (ii) $m = 2^\alpha$, $\alpha > 3$, $\lambda \equiv 5 \pmod{8}$, $r = 0$, we have

$$\tau D_\tau^{(2)} < \frac{1}{3} \sum_{i=0}^{k} |b_i| + k + \frac{17}{3},$$  \hspace{1cm} (11.13)

where $b_0, b_1, \ldots, b_k$ are the partial quotients in the continued fraction expansion to nearest integers of $\lambda/m$ in case (i) and of $\lambda/2^\alpha - 2$ in case (ii).

A slightly simpler bound is obtained if the discrepancy is only extended over dyadic squares contained in $I^2$ rather than over all subintervals of $I^2$ (cf. [57, Theorem 5.1]). This latter bound has been tabulated in [55], whereas values of local deviations relative to dyadic squares can be found in [57], [58]. The estimate (11.13) shows again that small partial quotients are the trademark of favorable multipliers. Case (ii) above is also in perfect harmony with (11.10), with both results stressing that in this case the arithmetic properties of $\lambda$ are more relevant modulo $2^\alpha - 2$ than modulo $2^\alpha$.

We shall demonstrate now that a result of the same type as Theorem 11.9 can be established without the technical apparatus of generalized Dedekind sums. Our estimate is in terms of standard continued fractions and its proof uses only elementary facts about continued fractions. The following auxiliary result is basic.

11.10. LEMMA. Let $m > 2$ and $\lambda$ be integers with $\gcd(\lambda, m) = 1$ and let $\theta$ be a real number. Then the discrepancy $D_N$ of the finite sequence of fractional parts $(\lfloor \lambda n/m \rfloor + \theta)$, $n = 0, 1, \ldots, N - 1$, satisfies

$$ND_N < \sum_{i=1}^{q} a_i \text{ for } 1 < N < m$$  \hspace{1cm} (11.14)

and $mD_m = 1$, where $a_1, \ldots, a_q$ are the partial quotients in the continued fraction expansion of $\lambda/m$. We also have

$$ND_N < C(K) \log(N + 1) \text{ for } 1 < N < m,$$  \hspace{1cm} (11.15)

where $K = K(\lambda/m) = \max(a_1, \ldots, a_q)$ and $C(K) = 2/\log 2$ for $1 < K < 3$, $C(K) = (K + 1)/\log(K + 1)$ for $K \geq 4$.

PROOF. If $N = m$, then the numbers $(\lfloor \lambda n/m \rfloor + \theta)$, $0 \leq n \leq m - 1$, form a sequence of $m$ equidistant points with distance $1/m$, and therefore we get $D_m = 1/m$. In the general case, we adapt a method from [220]. Let $1 = r_0 < r_1 < r_2 < \cdots < r_q = m$ be the denominators of the convergents to $\lambda/m$. Then an integer $N$ with $1 < N < m$ can be represented in the form $N = \sum_{i=0}^{h} c_ir_i$ with $0 < h < q$, $c_h \geq 1$, and $0 < c_i < a_{i+1}$ for $0 < i < h$. We decompose the given sequence $(\lfloor \lambda n/m \rfloor + \theta)$, $n = 0, 1, \ldots, N - 1$, into blocks, namely $c_h$ blocks of length $r_h$, $c_{h-1}$ blocks of length $r_{h-1}$, and so on. Consider such a block of length $r_i$, $0 < i < h$; it consists of elements of the form $(\lfloor \lambda n/m \rfloor + \theta)$, $n = n_0 + 1, \ldots, n_0 + r_i$. Let $p_i/r_i = [a_0; a_1, \ldots, a_i]$ be the $i$th convergent to $\lambda/m$. Then from the theory of continued fractions (cf. [241, p. 37]) we know that
Writing $n = n_0 + j$ with $1 \leq j < r$, we get

$$\left\{ \frac{\lambda n}{m} + \theta \right\} = \left\{ \frac{\lambda n_0}{m} + \frac{j p_i}{r_i} + \frac{j \delta_i}{r_i r_{i+1}} + \theta \right\} = \left\{ \frac{j p_i}{r_i} + u + \frac{j \delta_i}{r_i r_{i+1}} \right\},$$

where $u = (\lambda n_0 / m) + \theta$. Since $\gcd(p_i, r_i) = 1$, the numbers $(j p_i / r_i) + u$, $1 \leq j < r$, considered modulo 1, form a sequence of $r_i$ equidistant points with distance $1/r_i$, which therefore has discrepancy $1/r_i$. Because of $|j \delta_i / r_i r_{i+1}| < 1/r_{i+1}$ for $1 \leq j < r$, the finite sequence $\{ (\lambda n / m) + \theta \}$, $n = n_0 + 1, \ldots, n_0 + r_i$, is obtained by cyclically shifting modulo 1 the elements $\{(j p_i / r_i) + u\}$, $1 \leq j < r$, either all to the right or all to the left by at most a distance $1/r_{i+1}$, where the direction of the shift depends only on the sign of $\delta_i$. It is then easily seen that the discrepancy $D_{n_i}$ of the finite sequence $\{ (\lambda n / m) + \theta \}$, $n = n_0 + 1, \ldots, n_0 + r_i$, satisfies

$$D_{n_i} < 1/r_i + 1/r_{i+1}.$$

From the triangle inequality for discrepancies (cf. [174, p. 115]) and the way in which we decomposed the original sequence, it follows that the discrepancy $D_N$ of the sequence $\{ (\lambda n / m) + \theta \}$, $n = 0, 1, \ldots, N - 1$, satisfies

$$N D_N \leq \sum_{i=0}^{h} c_i \left( 1 + \frac{r_i}{r_{i+1}} \right) \leq \sum_{i=0}^{h} c_i \left( a_{i+1} / r_{i+1} \right) \leq \sum_{i=0}^{h} (c_i + 1),$$

where we used $c_i \leq a_{i+1}$ and $a_{i+1} r_i \leq r_{i+1}$ (cf. [241, p. 24]). The coefficients $c_i$ obtained from the algorithm in [220, p. 148] have the properties that $c_0 < a_1$ and that $c_i = a_{i+1}$ implies $c_{i-1} = 0$, and so the estimate (11.14) follows.

To obtain (11.15), it suffices to combine (11.16) with the inequality

$$\sigma(N) = \sum_{i=0}^{h} (c_i + 1) < C(K) \log(N + 1) \quad \text{for} \quad 1 \leq N < m,$$

where $\sigma(N)$ is well defined if we use the coefficients $c_i$ produced by the algorithm mentioned above. We establish (11.17) by induction on $h$. If $r_0 < r_1$, then the smallest possible $h$ is $h = 0$, and a corresponding $N$ satisfies $1 \leq N < r_1 < K$. If $r_0 = r_1 = 1$, then the smallest possible $h$ is $h = 1$ and a corresponding $N$ satisfies $1 \leq N < r_2 - 1 < K$. Since $\sigma(N) = N + 1$ for these $N$, it suffices to show for the first step in the induction that

$$N + 1 < C(K) \log(N + 1) \quad \text{for} \quad 1 \leq N < K.$$

The first property follows from $r_1 = a_1$, and for the second property we note that if $r_1 < N < r_{i+1}$, then $N - a_i r_i = N - a_{i+1} r_i < r_{i+1} - a_{i+1} r_i = r_{i-1}$. 

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But this follows from the fact that

\[ C(K) = \max_{1 < x < K} \frac{x + 1}{\log(x + 1)}. \]

Now take an arbitrary \( N \) with \( 1 < r_h < N < r_{h+1} \) and write \( N = c_h r_h + N_{h-1} \) with \( 0 < N_{h-1} < r_h \). Then \( \sigma(N) = c_h + 1 + \sigma(N_{h-1}) \), and the induction hypothesis yields \( \sigma(N) < c_h + 1 + C(K) \log(N_{h-1} + 1) \), which holds also for \( N_{h-1} = 0 \). Now \( N + 1 > c_h(N_{h-1} + 1) + N_{h-1} + 1 = (c_h + 1)(N_{h-1} + 1) \), and so

\[ \sigma(N) < c_h + 1 + C(K) \log \frac{N + 1}{c_h + 1}. \]

Since \( 1 < c_h < a_{h+1} < K \), we can complete the argument by applying (11.18) with \( c_h \) in place of \( N \).

11.11. Theorem. For a sequence of LCPRN with modulus \( m \), multiplier \( \lambda \), and either (i) \( r = m \); or (ii) \( m = 2^n, \alpha > 3, \lambda \equiv 5 \pmod{8}, r = 0 \), we have

\[ \tau D_r^{(2)} < 1 + \sum_{i=1}^{q} a_i \quad (11.19) \]

and

\[ \tau D_r^{(2)} < 1 + C(K) \log \tau, \quad (11.20) \]

where \( a_1, \ldots, a_q \) are the partial quotients in the continued fraction expansion of \( \lambda/m \) in case (i) and of \( \lambda/2^{\alpha-2} \) in case (ii), \( K = \max(a_1, \ldots, a_q) \), and \( C(K) \) is as in Lemma 11.10.

Proof. In case (i), all residues modulo \( m \) are generated by (7.1). Therefore the points \( x_0, x_1, \ldots, x_{m-1} \) are a permutation of the points

\[ \left\{ \frac{(n + r)}{m} \right\}, \quad n = 0, 1, \ldots, m - 1. \]

For a subinterval \( J = [u_1, u_2] \times [v_1, v_2] \) of \( I^2 \), the number \( A(J) \) of these points falling into \( J \) is equal to the number of integers \( n \) with \( mu_1 < n < mu_2 \) and \( \{\lambda n + r)/m \} \in [v_1, v_2] \). If the first condition is not satisfied by any integer \( n \), then

\[ |A(J) - m|J| = m|J| < 1. \quad (11.21) \]

Otherwise, the first condition is satisfied by the integers \( n = N_1, N_1 + 1, \ldots, N_2 - 1 \), say, where \( 0 < N_1 < N_2 \leq m \). Thus \( A(J) \) is equal to the number of integers \( j, 0 < j < N_2 - N_1 \), with

\[ \left\{ \frac{\lambda(N_1 + j) + r}{m} \right\} = \left\{ \frac{\lambda j}{m} + \theta \right\} \in [v_1, v_2], \]

where \( \theta = (\lambda N_1 + r)/m \). Hence we have \( A(J) = A([v_1, v_2); N_2 - N_1] \), with the second counting function referring to the sequence \( \{\lambda j/m + \theta\} \), \( j = 0, 1, \ldots, N_2 - N_1 - 1 \). If \( D_{N_2-N_1} \) is the discrepancy of this sequence, then

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|A(J) - m|J| | = |A([v_1, v_2]; N_2 - N_1) - m(u_2 - u_1)(v_2 - v_1)|
\leq |A([v_1, v_2]; N_2 - N_1) - (N_2 - N_1)(v_2 - v_1)|
+ |N_2 - N_1 - m(u_2 - u_1)|(v_2 - v_1)
\leq (N_2 - N_1)D_{N_2,N_1} + |N_2 - N_1 - m(u_2 - u_1)|.

Now \(N_i = mu_i + \theta_i\), \(0 < \theta_i < 1\), for \(i = 1, 2\), and so

\[|A(J) - m|J| | \leq (N_2 - N_1)D_{N_2,N_1} + 1 \leq \sum_{i=1}^{q} a_i + 1\]

by (11.14). Because of (11.21), this inequality holds for any \(J\), so that we obtain (11.19) for the case (i). The inequality (11.20) for the same case is shown by applying (11.15) instead of (11.14).

In case (ii), the recursion (7.1) generates all residues modulo \(m\) that are \(\equiv g \pmod{4}\), where \(g\) is the least residue of \(y_0\) modulo 4. Therefore the points \(x_0, x_1, \ldots, x_{r-1}\) are a permutation of the points

\([(4n + g)/m, \{\lambda(4n + g)/m\})\], \(n = 0, 1, \ldots, r - 1,\)

with \(r = m/4\). By using the same method as above, we obtain the desired results for case (ii).

The method in the proof of Theorem 11.11 can clearly be applied whenever the set of residues modulo \(m\) generated by (7.1) is known explicitly and consists of an arithmetic progression or possibly a union of arithmetic progressions. To provide one more example, consider the homogeneous case with a prime modulus \(m\) and a multiplier \(\lambda\) which is a primitive root modulo \(m\). Then \(r = m - 1\), and we get

\[(m - 1)D_{m-1}^{(2)} \leq 2 + \sum_{i=1}^{q} a_i\]

and

\[(m - 1)D_{m-1}^{(2)} \leq 2 + C(K) \log m\]

with the notation from case (i) in Theorem 11.11.

At first glance, the bound in (11.13) appears to be better than those in Theorem 11.11 by a factor \(\frac{1}{3}\). This superficial comparison ignores the fact that the \(|b_i|\) are usually larger than the \(a_i\). If we use, for instance, \(m = 2^{32}\), \(\lambda = 1812433253\), and \(r\) odd, then \(K = K(\lambda/m) = 2\), so that the bound in (11.20) becomes 65, whereas the bound in (11.19) is 50 and the bound in (11.13) is \(49\frac{1}{3}\).

Borosh and Niederreiter [24] have carried out a systematic search for those multipliers \(\lambda\) yielding a small value of \(K = K(\lambda/m)\) with respect to moduli \(m\) used in practice. The results show, for instance, that for each \(m = 2^a\), \(6 \leq \alpha \leq 35\), there exists a multiplier \(\lambda\) with \(K(\lambda/m) \leq 3\), and even one for which \(\lambda \equiv 5 \pmod{8}\). In many cases we can actually obtain \(K(\lambda/m) = 2\). An example with \(m = 2^{32}\) was already given above, and another interesting example is \(m = 2^{30}\), \(\lambda = 657759677\). Such multipliers yield, of course, a huge figure of merit \(\rho^{(2)}(\lambda, m)\) and an absolutely fantastic distribution of pairs.
Further results for the case \( s = 2 \) can be found in [229]. For instance, for every modulus \( m = 2^\alpha \) with \( \alpha > 31 \) there exists a \( \lambda \equiv 5 \pmod{8} \) with \( \rho(2)(\lambda, m) > m/\log m \), where \( m = 2m/5 \). The calculations mentioned above suggest that one may, in fact, achieve a lower bound of the order of magnitude \( m \). A similar existence theorem holds for odd prime power moduli \( m \) and multipliers \( \lambda \) that are primitive roots modulo \( m \). For a positive divisor \( d \) of an arbitrary modulus \( m \), the expression \( R^{(2)}(\lambda, m, d) \) introduced in Definition 11.1 can be estimated in terms of \( K(\lambda/d) \).

With all this information for \( s = 2 \), it can be said that the distribution of pairs of successive LCPRN is now clearly understood. For the moduli of practical interest, the discrepancy over the full period \( x_0, x_1, \ldots, x_{T-1} \) is \( O(\tau^{-1} \log \tau) \) with suitable multipliers that can be determined effectively (cf. [24]). In the light of (3.11), this order of magnitude is the lowest possible for any \( \tau \) points in \( I^2 \), so that well-chosen LCPRN lead to a nearly optimal distribution of pairs. In order to pass the 2-dimensional serial test, the decisive factor in the proper selection of a multiplier is the continued fraction expansion of \( \lambda/m \) (or of a closely related rational) which should only involve small partial quotients.

Proposals for the choice of multipliers have been based on weaker criteria, notably those of small serial correlation. The estimate of Greenberger [86] for the serial correlation suggested to take \( \lambda \approx m^{1/2} \), but this yields \( \rho(2)(\lambda, m) = O(m^{1/2}) \) and thus a distribution of pairs of modest quality. The same objection can be raised against any multiplier that is too small in comparison with the modulus. Even a generator such as \( m = 2^{35}, \lambda = 2^{24} + 5, r = 1 \), for which Jansson [141] reported excellent serial correlation properties, is not very impressive in terms of its distribution of pairs since we have \( \rho(2)(\lambda, m) < \lambda \approx m^{2/3} \). It must also be pointed out that multipliers which are too close to a fairly large power of 2 may permit a fast generation procedure, but perform very poorly under the 3-dimensional serial test. For instance, Jansson's generator satisfies \( \lambda^2 - 10\lambda + 25 \equiv 0 \pmod{m} \), and so \( \rho(3)(\lambda, m) < 250 \), which is much too small.40 Another generator of this type that has been used in some computer installations is \( m = 2^{35}, \lambda = 2^{18} + 3, r = 0 \). Here we have \( \gamma = 3 \) and \( \lambda^2 - 6\lambda + 9 \equiv 0 \pmod{2^{2\gamma-r}} \), so that \( \rho(3)(\lambda, 2^{2\gamma-r}) < 54 \), which is totally unacceptable. The fact that such multipliers yield a bad distribution of triples was already noted earlier on the basis of numerical evidence (cf. [48], [87], [184], [342], [345]). The proposal of Ahrens, Dieter, and Grube [1] mentioned in §10 is thoroughly reasonable, although it was based on the weaker argument of small serial correlation and does not necessarily yield optimal multipliers with respect to the distribution of pairs.

Specific multipliers have also been proposed on the basis of the "cubic-lattice criterion" (see §10). Marsaglia [188, p. 275] comes out strongly in favor of \( m = 2^{32}, \lambda = 69069 \), which yields an almost cubic 2-dimensional lattice in the sense that it has a unit cell for which the longer side is about 1.06 times the shorter side. However, this does not imply a good distribution of pairs. Concretely, we have \( 69069 \cdot 1 - 1 \cdot \lambda \equiv 0 \pmod{m} \), and so \( \rho(2)(\lambda, m) < \)

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40 This is an illustration of the fact that "going up" in the dimension may change the performance of a multiplier completely.
Thus, if we use this multiplier in the mixed congruential method with \( r \) odd, it follows from Theorem 11.8 with \( C'_2 = 1/4 \) that \( mD^{(2)}_m > 15545 \). On the other hand, \( m = 2^{32} \), \( \lambda = 1589013525 \), and \( r \) odd yields \( mD^{(2)}_m < 48 \) by (11.13).\(^{41}\) These data demonstrate more eloquently than anything else the inutility of the “cubic-lattice criterion”. In fact, one would be better off choosing multipliers blindly rather than using this criterion, for Yao and Knuth [364] proved that the expected value of the bound in (11.19) is \((6/\pi^2)\log^2 m\) plus lower-order terms, which still leads to a satisfactory value for \( D^{(2)}_m \).

We recapitulate that the performance of a sequence of LCPRN under the \( s \)-dimensional serial test is governed by a certain \( s \)-dimensional figure of merit which must be large for a good generator. We should observe the relativity principle and the general rule that a multiplier not be related to the modulus in too simple a manner. A saving of time in the generation procedure may be punished by a serious distortion in the numerical work for which the LCPRN are used. If the parameters are chosen according to the criteria spelled out in this section, then the \( s \)-tuples of successive LCPRN show an excellent distribution behavior in \( I^s \), and so we have statistical almost-independence among \( s \) successive terms. In view of the simplicity of the recursion (7.1), it is quite astonishing how close the distribution of \( s \)-tuples can be to an optimal one. Marsaglia’s devastating verdict that “congruential random number generators are not suitable for precision Monte Carlo use” [188, p. 250] can be overruled on the basis of new evidence.

Even among supporters of Lehmer’s method one may occasionally sense a lingering suspicion about statistical independence properties (cf. [154, p. 56]). The standard argument leading to such misgivings has it that (7.1) establishes such an intimate link between successive terms that it is futile to hope for independence among these terms. In order to understand why LCPRN nevertheless perform so well, the following analogy may be helpful. In parentheses, we provide the translation from the model to LCPRN. Consider a square pool table and a billiard ball moving on it with enormously high speed (\( = \) the multiplier \( \lambda \) is very large). The tangent of the angle between a boundary of the table and the initial direction of the motion is supposed to be an irrational number (\( = \) for a good generator, the fraction \( \lambda/m \) has approximately the arithmetic properties of an irrational number). In one time unit, the billiard ball undergoes several thousand reflections at the boundaries (\( = y_{n+1} \) is obtained from \( \lambda y_n + r \) by subtracting a multiple of \( m \) that is usually very large). The path of the billiard ball is locally a straight line and thus predictable, but the position at time \( n + 1 \) is practically independent of that at time \( n \). The high velocity and the many reflections “smear out” any correlation that might exist locally (\( = y_{n+1} \) is almost independent of \( y_n \)). It is a known fact (cf. [174, p. 87, Exercise 9.29]) that under the given conditions the billiard ball visits, in the long run, every rectangular region and, more generally, every Jordan-measurable region on the table for an amount of time proportional to the area of that region (\( = \) a sequence of LCPRN with large

\(^{41}\) The simpler bound (11.19) shows \( mD^{(2)}_m < 51 \), whereas (11.20) leads to \( mD^{(2)}_m < 65 \) since we have \( K = K(\lambda/m) = 2 \).
period and a properly chosen multiplier has an excellent distribution behavior. Thus, the pseudo-randomness properties of a good sequence of LCPRN are comparable to those implied by the determinate, but nevertheless wildly erratic and for all practical purposes "pseudo-Brownian" motion of a "supersonic" billiard ball.

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