Research Article

An Efficient Differential MIMO-OFDM Scheme with Coordinate Interleaving

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We propose a concatenated trellis code (TC) and coordinate interleaved differential space-time block code (STBC) for OFDM. The coordinate interleaver, provides signal space diversity and improves the codeword error rate (CER) performance of the system in wideband channels. Coordinate interleaved differential space-time block codes are proposed and used in the concatenated scheme, TC design criteria are derived, and the CER performances of the proposed system are compared with existing concatenated TC and differential STBC. The comparison showed that the proposed scheme has superior diversity gain and improved CER performance.

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1. INTRODUCTION

In recent years, code design for multiple-input multiple-output (MIMO) channels, with orthogonal frequency division multiplexing (OFDM) modulation, has gained much attention in wireless communications. Space-time block codes (STBC) first proposed by Alamouti [1] provide full spatial diversity in wireless channels, with simple linear maximum likelihood (ML) decoder. An efficient scheme of concatenated trellis code and STBC (TC-STBC) which provides additional diversity and coding gain was proposed by Gong and Letaief [2]. Tarasak and Bhargava [3] applied the constant modulus (CM) differential encoding scheme of Tarokh and Jafarkhani [4] to the TC-STBC system [2]. The differential encoding has the advantage of avoiding channel estimation and the transmission of pilot symbols. Further improvement of TC-STBC performance is possible by using a coordinate interleaver [5]. Coordinate interleaved signal sets provide signal space diversity and hence improve the symbol error performance of communication systems in fast fading channels. The recent application of coordinate interleaving to MIMO-OFDM which shows that this technique provides considerable diversity gain without significant increase of encoding and decoding complexities was proposed by Rao et al. [6]. The single symbol decodability of coordinate interleaved orthogonal design (CIOD) [7] is an important feature ensuring low decoding complexity. The joint use of CIOD and OFDM provides spatial and multipath diversities, and further concatenation of TC and CIOD (TC-CIOD) [5] as a consequence gives much better performance compared to CIOD OFDM [6], linear constellation precoded (LCP)-CIOD OFDM [6], and TC-STBC OFDM [2].

In this paper, we apply the nonconstant modulus (non-CM) differential space-time block (STB) encoding scheme proposed by Hwang et al. [8] to CIOD, and use it in TC-CIOD scheme [5]. The proposed differential scheme achieves full spatial and multipath diversities, and provides considerable coding gain advantage without channel state information (CSI). We derive the design criteria for differential TC-CIOD and found that under some approximation they are same as in TC-CIOD case. The new differential scheme provides same diversity gain as the TC-CIOD scheme, and has diversity four times greater than of both the TC-STBC system introduced by Gong and Letaief [2] and its differential counterpart proposed by Tarasak and Bhargava [3]. To clarify the effect of interleaver selection on the diversity gain of TC-STBC, we extend the results given in [2, 3] where the
two-symbol interleaver is considered between TC and STBC, to the symbol interleaver case.

2. PRELIMINARIES

In this section, we summarize the encoding and decoding of non-CM differential STBC, and in the following sections, the non-CM differential STBC is used in differential TC-CIOD system. Note that, the use of any CM differential encoding technique with CIOD is not possible due to nonconstant modulus of coordinate interleaved signal constellation.

Let us assume a quasistatic fading channel with two transmit and one receive antennas, and denote the channel gains corresponding to two transmit antennas with \( h_1 \) and \( h_2 \), respectively. Let the dummy symbols to be transmitted during the first two transmission periods be \( a_1 \) and \( a_2 \). Therefore, \( a_1 \) and \( -a_2 \) are transmitted from the first transmit antenna, and \( a_2 \) and \( a_1^* \) are transmitted from the second transmit antenna during the first and second transmission periods, respectively. The differential STBC encodes the first data symbol pair \((x_1, x_2)\) by using the following equations [8]:

\[
\begin{align*}
a_1 &= \frac{x_1a_1 - x_2a_2^*}{\sqrt{|a_1|^2 + |a_2|^2}}, \\
a_4 &= \frac{x_1a_2 + x_2a_1^*}{\sqrt{|a_1|^2 + |a_2|^2}}.
\end{align*}
\]

(1)

The difference of non-CM differential STBC from CM differential STBC [4] is in the scaling coefficient \( \sqrt{|a_1|^2 + |a_2|^2} \) which ensures that the total transmission energy of two antennas remains equal to one. The transmission of space-time-block- (STB-) encoded dummy symbols \( a_1 \) and \( a_2 \) results with reception of

\[
\begin{align*}
r_1 &= h_1a_1 + h_2a_2 + n_1, \\
r_2 &= -h_1a_2^* + h_2a_1^* + n_2,
\end{align*}
\]

(2)

where \( n_1 \) and \( n_2 \) are complex additive white Gaussian noise terms. Similarly, the transmission of STB-encoded \( a_3 \) and \( a_4 \) carrying non-CM symbols \( x_1 \) and \( x_2 \) results with reception of

\[
\begin{align*}
r_3 &= h_1a_3 + h_2a_4 + n_3, \\
r_4 &= -h_1a_4^* + h_2a_3^* + n_4.
\end{align*}
\]

(3)

The differential decoder uses the received symbols \( r_1, r_2, r_3, \) and \( r_4 \) to find the estimations of the transmitted non-CM symbols using

\[
\begin{align*}
\hat{x}_1 &= \frac{r_1r_4^* + r_2r_3}{(|h_1|^2 + |h_2|^2)|a_1|^2 + |a_2|^2}, \\
\hat{x}_2 &= \frac{r_2r_4^* - r_1r_3}{(|h_1|^2 + |h_2|^2)|a_1|^2 + |a_2|^2}.
\end{align*}
\]

(4)

As seen from (4), to find the transmitted non-CM symbol estimates \( \hat{x}_1 \) and \( \hat{x}_2 \), the receiver should know or at least estimate the channel power \( (|h_1|^2 + |h_2|^2) \) and the signal power of previously transmitted symbols \( (|a_1|^2 + |a_2|^2) \).

The simple estimation for the channel power \( p = (|h_1|^2 + |h_2|^2) \) denoted by \( \hat{p} \) is possible by evaluating the expected value of \( |r_i|^2 \), that is,

\[
\hat{p} = \frac{M|R|^H}{M},
\]

(5)

where \( M \) is the number of received symbols included in expected value calculation, \( R = [r_1 \ r_2 \ r_3 \ r_4 \ \cdots \ r_M] \), and \( R^H \) is the Hermitian of \( R \). The computational complexity of (5) can be reduced by using

\[
\hat{p}' = M - 1 - \frac{1}{M}|r_t|^2 + \frac{1}{M}|r_t|^2,
\]

(6)

where \( t \) is the recursion index.

There are two simple methods to estimate the signal power of previously transmitted symbols. The first one is to use the previous decoder output. The second one is to use (2) to obtain

\[
|r_1|^2 + |r_2|^2 = (|h_1|^2 + |h_2|^2)(|a_1|^2 + |a_2|^2) + n_r,
\]

(7)

where \( n_r \) is the Gaussian noise term. From (7), the estimation of the signal power of previously transmitted symbols can be written as

\[
(|a_1|^2 + |a_2|^2) \approx \frac{|r_1|^2 + |r_2|^2}{\hat{p}}.
\]

(8)

3. SYSTEM MODEL

In this section, we describe the proposed differential TC-CIOD OFDM system, and its encoding and decoding operations.

3.1. Differential encoder

The encoder block diagram of the proposed differential TC-CIOD OFDM for two transmit antennas is shown in Figure 1, where the source bits are trellis encoded at rate 2/3 and mapped to 8-PSK signal constellation. Each 8-PSK symbol is rotated by \( \theta \) and then a vector of rotated symbols is coordinate interleaved by \( \pi \). To achieve maximum diversity, a proper coordinate interleaver should be used. Let

\[
\bar{X}^t = [x_0^t \ x_1^t \ x_2^t \ x_3^t \ \cdots \ x_{K-2}^t \ x_{K-1}^t]
\]

(9)

be the \( t \)th rotated trellis codeword of length \( 2K \), where the symbols \( \bar{x}_k^t \) are obtained by rotating the symbols \( x_k^t \) of the \( t \)th trellis codeword \( X^t \) by \( \theta \), that is,

\[
\bar{x}_k^t = x_k^t \exp (j\theta).
\]

(10)

The coordinate interleaver \( \pi_t \), which has a great impact on the overall system performance, performs the following assignments:

\[
\begin{align*}
\bar{x}_{2k}^t &= x_{k,J}^t + jx_{k,K/2+J}^t, \\
\bar{x}_{2k+1}^t &= x_{K/2+k,J}^t + jx_{K+J+k,(3K/2)}^t.
\end{align*}
\]

(11)
for $k = 0, \ldots, K - 1$, and the coordinate interleaved symbols $\tilde{X}_k^\nu$ form the vector

$$\tilde{X}_k^\nu = \left[ X_{k0}^\nu, X_{k1}^\nu, X_{k2}^\nu, \ldots, X_{K-2}^\nu, X_{K-1}^\nu \right].$$  \hfill (12)

In (11), the operators $(\cdot)_t$ and $(\cdot)_Q$ represent the real and imaginary parts of a complex symbol, respectively, and the operator $(\cdot)_2$ takes modulo $2K$ of the operand. The vector $\tilde{X}_k^\nu$ enters the differential encoder which produces a vector $A^{t+1}$ with elements $\alpha_{k}^{t+1}$ obtained from

$$\begin{align*}
\alpha_{2k}^{t+1} &= \frac{X_{2k}^{t}d_{2k}^{t} - X_{2k+1}^{t}d_{2k+1}^{t}}{\sqrt{|d_{2k}^{t}|^2 + |d_{2k+1}^{t}|^2}}, \\
\alpha_{2k+1}^{t+1} &= \frac{X_{2k}^{t}d_{2k+1}^{t} + X_{2k+1}^{t}d_{2k}^{t}}{\sqrt{|d_{2k}^{t}|^2 + |d_{2k+1}^{t}|^2}}
\end{align*}$$  \hfill (13)

for $k = 0, \ldots, K - 1$, similar to (1). The differentially encoded symbol pairs $\alpha_{2k}^{t+1}$ and $\alpha_{2k+1}^{t+1}$ are STB encoded as

$$Y_k^{t+1} = \begin{pmatrix} \alpha_{2k}^{t+1} \\ \alpha_{2k+1}^{t+1} \\ \alpha_{2k}^{t+1*} \\ \alpha_{2k+1}^{t+1*} \end{pmatrix}$$  \hfill (14)

and transmitted from the $\alpha_k$th OFDM subcarrier. There is a one-to-one mapping between $k$ and OFDM subcarriers, denoted by $\alpha_k$, which corresponds to the channel interleaver $\alpha$. The rows of $Y_k^{t+1}$ are transmitted from $(2t + 2)$th and $(2t + 3)$th OFDM frames, respectively, and the columns of $Y_k^{t+1}$ are transmitted from first and second transmit antennas, respectively.

The differential transmitter starts encoding at $t = 0$ by using initial dummy vector $A^0$ with nonzero elements selected from considered signal constellation. The transmission consists of first STB encoding of arbitrary vector $A^0$, which does not convey any information, and then sending it in the first two OFDM frames. The transmitter subsequently encodes the data in an inductive manner.

### 3.2. Channel model

Multipaths between transmit and receive antenna pairs in wireless communication channels cause intersymbol interference (ISI) in the received signals. The baseband impulse response for the MIMO channel with $L$ paths between the $\mu$th transmit ($1 \leq \mu \leq n_T$) and $v$th receive ($1 \leq v \leq n_R$) antennas is given as [9]

$$h_{\mu v}(t, \tau) = \sum_{l=0}^{L-1} h_{\mu v}(t, l) \delta(\tau - \tau_l).$$  \hfill (15)

In (15) $h_{\mu v}(t, l)$ is the time-dependent channel tap weight, $\delta(\cdot)$ is the Dirac function, and $\tau_l$ is the path propagation delay of the $l$th path ($0 \leq l \leq L - 1$). OFDM modulation with cyclic prefix (CP) addition at the transmitter and removal at the receiver transforms the frequency-selective channel into $K$ frequency nonselective subchannels without ISI. Assuming that the channel weights remain constant during an OFDM frame, the channel response becomes independent from time variable $t$, for single OFDM symbol period, and then the signal received by the $v$th antenna at the $t$th symbol interval, for the $k$th subcarrier ($0 \leq k \leq K - 1$), can be expressed as

$$r_{k}^v(t) = \sum_{\mu=1}^{n_T} H_{\mu v}^k(t) y_{\mu v}^k(t) + n_{k}^v(t),$$  \hfill (16)

where $y_{\mu v}^k(t)$ is the symbol transmitted by the $k$th subcarrier during $t$th symbol interval from $\mu$th transmit antenna, the samples $n_{k}^v(t)$ are zero-mean complex Gaussian r.v. with variance $\sigma_0^2$ per dimension, and

$$H_{\mu v}^k(t) = \sum_{l=0}^{L-1} H_{\mu v}^k(t) \exp \left(-\frac{j2\pi k\tau_l}{T_s} \right)$$  \hfill (17)

is the frequency-domain complex subchannel gain between $\mu$th transmit and $v$th receive antennas for the $k$th subchannel during $t$th symbol interval. In (17), $T_s$ is the effective OFDM symbol interval length and $h_{\mu v}^k(0)$ is the channel tap weight.

For simplicity we will drop the receive antenna index $v$ in the following derivations. However, the proposed system structure is easily extensible for more than one receive antenna. If we assume a quasistatic channel, we may also drop the time index $t$, from subcarrier transmission gains. Let the transmission of $Y_k^{t+1}$ be affected by the subcarrier transmission gains $H_k^k(\alpha_k)$ and $H_v^k(\alpha_k)$ corresponding to the first and second transmit antennas, respectively. For simplicity, we will denote $H_{\mu v}(\alpha_k)$ as $H_{\mu v}^k$, for $\mu = 1, 2$. Let, $r_{k}^v$ be the symbol received from $\alpha_v$th subcarrier of the $(t + 1)$th OFDM symbol, and $n_{k}^v$ for $k = 0, 1, \ldots, K - 1$ being the subchannel noise variables which are independent and identically distributed zero-mean complex Gaussian r.v. with variance $\sigma_0^2$ per dimension. Then, the MIMO-OFDM transmission can be modeled by

$$R_k^{t+1} = Y_k^{t+1} H_k + N_k^{t+1},$$  \hfill (18)

where $R_k^{t+1} = \{ r_k^{2t+2}, r_k^{2t+3} \}$, $H = \{ H_k^0, H_k^1 \}$, and $N_k^{t+1} = \{ n_k^{2t+2}, n_k^{2t+3} \}$ for $k = 0, 1, \ldots, K - 1$. Let us consider an OFDM codeword $Y_k^{t+1} = \{ Y_0^{t+1}, Y_1^{t+1}, Y_2^{t+1}, \ldots, Y_{K-1}^{t+1} \}$ transmuted over $K$ different subcarriers. The transmission of codeword $Y_k^{t+1}$ results in the reception of $R_k^{t+1} = \{ R_0^{t+1}, R_1^{t+1}, R_2^{t+1}, \ldots, R_{K-1}^{t+1} \}$, and the corresponding additive Gaussian noise affecting $R_k^{t+1}$ can be expressed as $N_k^{t+1} = \{ N_0^{t+1}, N_1^{t+1}, N_2^{t+1}, \ldots, N_{K-1}^{t+1} \}$.

### 3.3. Differential decoder

When the receiver does not have any CSI, the decoding metric for the trellis codeword

$$X^t = \left[ x_0^t, x_1^t, x_2^t, \ldots, x_{K-2}^t, x_{K-1}^t \right]$$  \hfill (19)
can be expressed as
\[
m(R^{t+1}, R', X') = \sum_{k=0}^{(K/2)-1} m_k'.
\]
(20)

The decoder should determine the \(t\)th trellis codeword \(X_t\) minimizing (20) to perform maximum likelihood (ML) decoding, where the different CIOD decoding metric is defined as
\[
m_k' = m\left[R_k^{t+1}, R_{k+(K/2)}^{t+1}, R_k^{t+1}, R_{k+(K/2)}^{t+1}, x_{2k}, x_{2k+1}, x_{2k+K}, x_{2k+K+1}\right].
\]
(21)

The CIOD decoding metric \(m_k'\) used in (20) can be written as
\[
m_k' = \tilde{m}_k' + \tilde{m}_{2k+1} + \tilde{m}_{2k+K} + \tilde{m}_{2k+K+1}'
\]
(22)
for \(k = 0, \ldots, (K/2) - 1\), where the STB symbol metric for \(\xi = 2k, 2k + 1, 2k + K\) and \(2k + K + 1\) is
\[
\tilde{m}_k' = |\hat{x}_k' - \hat{x}_k|^2 + (S_k' - 1)|\hat{x}_k|^2,
\]
(23)
derived similar to [10, page 453]. In (23), the scaling coefficient, which can be estimated by the methods described at the end of Section 2, is given by
\[
S_k' = (|H_k|^2 + |H_k|^2)\sqrt{|a_{2k}|^2 + |a_{2k+1}|^2}
\]
(24)
and the coordinate interleaved symbol estimates for \(k = 0, \ldots, K - 1\) are
\[
\hat{x}_{2k} = r_{2k}^2 + r_{2k+1}^2 + r_{2k+3}^2 + r_{2k+4}^2,
\]
\[
\hat{x}_{2k+1} = r_{2k}^2 + r_{2k+1}^2 + r_{2k+3}^2 + r_{2k+4}^2
\]
(25)
similar to (4). The scaling coefficient in (24) can be estimated by using the subchannel power estimation as (6) and the signal power estimation of previously transmitted symbols as (8). Similar to (6) and (8), we can express the estimation of \(S_k'\) as
\[
\hat{S}_k' = \sqrt{\hat{p}_k\left(|r_{2k}^2|^2 + |r_{2k+1}^2|^2\right)}
\]
(26)
where the subchannel power estimate \(\hat{p}_k\) is calculated recursively from
\[
\hat{p}_k' = \frac{M - 2}{M} \hat{p}_k' - 2 \frac{M}{M} \left(|r_{2k}^2|^2 + |r_{2k+1}^2|^2\right)
\]
(27)
with initial value \(\hat{p}_0' = 1\).

The metrics in (22) related with coordinate interleaved STB symbols are not suitable for Viterbi decoding. Substituting (23) in (22) and using (11), the metrics in (22) become related to the rotated trelliss codewords symbols \(\bar{x}_k, x_{k+(K/2)}, x_{k+K}\), and \(x_{k+(3K/2)}\). Hence, the CIOD decoding metric in (22) can be further expressed in terms of the branch metrics \(m_k^t\), \(m_{2k+1}^t\), \(m_{2k+K}^t\), and \(m_{2k+K+1}^t\) as
\[
m_k' = m_k^t + m_{2k+1}^t + m_{2k+K}^t + m_{2k+K+1}^t.
\]
(28)

where
\[
m_k^t = (x_{2k,i} - \bar{x}_{2k,i})^2 + (x_{2k+1,i} - \bar{x}_{2k+1,i})^2 + (x_{2k+K,i} - \bar{x}_{2k+K,i})^2 + (x_{2k+K+1,i} - \bar{x}_{2k+K+1,i})^2
\]
\[
m_{2k+1}^t = (x_{2k+1,i} - \bar{x}_{2k+1,i})^2 + (x_{2k+K+1,i} - \bar{x}_{2k+K+1,i})^2
\]
\[
m_{2k+K}^t = (x_{2k+K,i} - \bar{x}_{2k+K,i})^2 + (x_{2k+1,i} - \bar{x}_{2k+1,i})^2
\]
\[
m_{2k+K+1}^t = (x_{2k+K+1,i} - \bar{x}_{2k+K+1,i})^2 + (x_{2k,i} - \bar{x}_{2k,i})^2
\]
(29)
for \(k = 0, \ldots, (K/2) - 1\), which can be used by Viterbi decoder, to estimate the source bits.

4. TRELLIS CODE DESIGN

To achieve full diversity and high coding gain with the proposed differential TC-CIOD OFDM, we obtained the pairwise error probability (PEP) upper bound, which is the probability that the decoder chooses an erroneous sequence \(Z\) instead of the transmitted sequence \(X\), defined as
\[
P(X, Z | H) = \Pr\left[m(R^{t+1}, R', X') > m(R^{t+1}, R', Z')\right]
\]
(30)

In (30), we substitute \(m(R^{t+1}, R', X')\) with the metrics in (20), (22), and (23), and the corresponding metrics for \(m(R^{t+1}, R', Z')\). Assuming that the previous codeword symbols \(a_k = (1 + j)/2\) and the subchannel noise variables \(n_k^t\) are i.i.d. zero-mean complex Gaussian distributed r.v. with variance \(N_0/2\) per dimension, by dropping the time index \(t\) for simplicity, we obtain
\[
P(X, Z | H) = \prod_{k=0}^{(K/2)-1} Q\left[\frac{E_t}{2N_0 d_k^2(1+E_t^2)+d_{k+(K/2)}^2(1+E_t^2)}\right]
\]
(31)
where \(Q(\cdot)\) is the Gaussian error function:
\[
d_k^2 = \sum_{i=0}^{1} (|H_{k,i}|^2 + |H_{k,i}|^2) |\hat{x}_{2k+i} - \bar{x}_{2k+i}|^2
\]
(32)
and the symbol energy involved in STBC is
\[
E_t^2 = \sum_{i=0}^{1} |\bar{x}_{2k+i}|^2
\]
(33)
for \(\xi = k \text{ and } k + (K/2)\). If we further assume that \(E_t^2 = 1\), the pairwise error probability given by (31) simplifies to
\[
P(X, Z | H) = \prod_{k=0}^{(K/2)-1} \sqrt{\frac{E_t}{4N_0 \left[d_k^2 + d_{k+(K/2)}^2\right]}}
\]
(34)
which is the same expression given in [5], except that 2N₀ is replaced by 4N₀, corresponding to 3 dB performance loss of differential TC-CIOD scheme. Using the inequality

$$Q(x) \leq \frac{1}{2} \exp \left( - \frac{x^2}{2} \right)$$

(35)

and ignoring multiplier 1/2 for simplicity, we may upper bound (34) as

$$P(X, Z | H) < \exp \left[ - \frac{E_s}{8N_0} d^2(X, Z) \right],$$

(36)

where the modified Euclidean distance between pair of trellis codewords X and Z is given as

$$d^2(X, Z) = \sum_{k=0}^{K-1} \sum_{i=0}^{1} \left( |H_{f,k}^i|^2 + |H_{g,k}^i|^2 \right) |\tilde{x}_{2k+i} - \tilde{z}_{2k+i}|^2.$$  (37)

The rotated trellis codewords corresponding to X and Z are denoted by X and Z, respectively. Let X and Z differ only during the short part with length κ, that is, only [x₁, x₂, ..., xₖ] differs from [z₁, z₂, ..., z₂ₖ]. In this case, we may rewrite (37) as

$$d^2(X, Z) = \sum_{k=1}^{s+\kappa} \left( |H_{f,k}^ι|^2 |\tilde{x}_{k} - \tilde{z}_{k}|^2 + |H_{g,k}^ι|^2 |\tilde{x}_{k,Q} - \tilde{z}_{k,Q}|^2 \right),$$

(38)

where ι = {s+1, s+2, ..., s+κ}, f(k) = ⌊π₁(k)/2⌋, g(k) = ⌊π₂(k)/2⌋ and ⌊·⌋ takes the integer part of the operand. The coordinate interleaver π can be represented by a pair of permutations for real and imaginary parts of the input vector denoted by π₁(k) and π₂(k), respectively, in the definition of f(k) and g(k). According to (11),

$$π₁(k) = \begin{cases} 2k, & k < K, \\ 2k - 2K + 1, & K \leq k < 2K, \end{cases}$$

(39)

$$π₂(k) = \begin{cases} 2k + K + 1, & k < \frac{K}{2}, \\ 2k - K, & \frac{K}{2} \leq k < \frac{3K}{2}, \\ 2k - 3K + 1, & \frac{3K}{2} \leq k < 2K. \end{cases}$$

(40)

Perfect coordinate interleaving guarantees that f(ξ) ≠ g(ω) for every pair ξ, ω ∈ ι and f(ξ) ≠ f(ω), g(ξ) ≠ g(ω) for every pair ξ, ω ∈ ι when ξ ≠ ω. Assuming perfect coordinate interleaving, there are no repeated subcarrier fading coefficients Hₖᵢ in (38). If the subcarriers are perfectly interleaved and transmit antennas are well separated, we can assume that the subcarrier fading coefficients Hₖᵢ used in (38) are zero mean i.i.d. complex Gaussian random variables with variance 1/2 per dimension. Taking the expectation of (36) over Rayleigh distributed r.v. |Hₖᵢ| using (38), we obtain

$$P(X, Z) < \prod_{k \in ι} \left[ 1 + \frac{E_s}{8N_0} (\tilde{x}_{k,J} - \tilde{z}_{k,J})^2 \right]^{-1} \left[ 1 + \frac{E_s}{8N_0} (\tilde{x}_{k,Q} - \tilde{z}_{k,Q})^2 \right]^{-1},$$

(41)

In general, θ can be selected such that for xₖ ≠ zₖ, both of real and imaginary components of xₖ and zₖ do not differ. Hence, we should consider two different sets of k values, ι and ι₂, for which real and imaginary components of rotated trellis codeword symbols xₖ and zₖ differ, respectively. In this case, at high signal-to-noise ratios (SNR), (41) can be expressed as

$$P(X, Z) < \left( \frac{E_s}{8N_0} \right)^{-2|ι|+|ι₂|} \left[ \prod_{k \in ι} (\tilde{x}_{k,J} - \tilde{z}_{k,J}) \prod_{k \in ι₂} (\tilde{x}_{k,Q} - \tilde{z}_{k,Q}) \right]^{-1},$$

(42)

where |ι| and |ι₂| represent the cardinality of sets ι and ι₂, respectively. It is clear from (42) that under the assumption of perfect coordinate and channel interleaving, the achievable diversity of the system is

$$G_d = 2 \times \min_{X, Z} (|ι| + |ι₂|),$$

(43)

and the differential TC-CIOD coding gain is

$$G_c = \frac{1}{2} \min_{X, Z} \left[ \prod_{k \in ι} (\tilde{x}_{k,J} - \tilde{z}_{k,J}) \prod_{k \in ι₂} (\tilde{x}_{k,Q} - \tilde{z}_{k,Q}) \right]^{-\frac{1}{2G_d}},$$

(44)

The codeword error probability can be written in terms of pairwise error probability as

$$P_e = \sum_{X} \sum_{Z \in X} P(X, Z),$$

(45)

where P(X) is the probability of the codeword X being generated by the trellis encoder and the PEP P(X, Z) is upper bounded by (42). The trellis code and θ can be selected to minimize the codeword error probability upper bound obtained by substituting (42) in (45). The trellis code search is performed over all possible trellis generator polynomials based on the representation given in [11]. We selected θ values ranging from 0.5° till 22.5° with 2° steps and E_s/N₀ = 17 dB during an exhaustive computer-based 4-, 8-, 16-, and 32-state 8-PSK R = 2/3 trellis codes search minimizing the codeword error probability upper bound calculated over all possible trellis codeword pair X and Z with length κ = 3 starting and ending at the common trellis states. Figure 2 shows the codeword error probability (P_e) upper bound of best trellis codes found for different values of θ for considered 4-, 8-, 16-, and 32-state trellises. It is clear from Figure 2 that the codeword error probability upper bounds for the best trellis code decrease with θ and achieve their minimum
In this section, we give the simulation results for the proposed system and evaluate the effect of interleaver selection on the performance of the concatenated schemes. We use two-symbol [3], symbol, and coordinate interleavers and consider the performance of both differential and nondifferential TC-STBCs. Figure 3 shows the codeword error rate (CER) of the systems with efficiency of 2 bps/Hz, when trellis code termination and OFDM cyclic prefix are excluded. The channel model used during the simulations is given in (18), where $H_k^i$s are independent and identically distributed Gaussian random variables with variance 1/2 per dimension, and in order to obtain the mean CER performances of the differential systems, the $H_k^i$ values are randomly assigned multiple times during the simulation after each 10 codeword transmissions followed by a dummy frame transmission to initiate the differential decoder to the random channel change. Hence, this model corresponds to a very slow varying fading channel. The perfectly interleaved multipath channel, that is, independent $H_k^i$s, 48 OFDM subcarriers, and the perfect knowledge of the scaling coefficients $S_k^i$s, were assumed during the simulations. The proposed scheme outperforms the differential two-symbol interleaved TC-STBC proposed by Tarasak and Bhargava [3] by 8.5 dB in SNR at a CER of $10^{-3}$. Note that the symbol interleaver doubles the multipath diversity achieved by TC-STBC compared to two-symbol interleaver considered in [2, 3], and outperforms the two-symbol interleaved case by 6.5 dB in SNR at the CER of $10^{-3}$. During the simulations, we employed a $2 \times 48$ block interleaver between TC and STBC as symbol interleaver. When a symbol interleaver is used, the set size $\omega_i$ defined in [2], becomes equal to effective length (time diversity) of the trellis code. Hence, the maximum achievable diversity of TC-STBC doubles. All of the codes employ a rate 2/3 8-PSK 4-state trellis used in [2], except the one denoted by T2, which uses the optimized 4-state trellis code given in Table 1. For TC-CIOD, the rotation angle $\theta$ is taken equal to 22.5°, which is found to be optimum for $R = 2/3$ 8-PSK trellis codes with 4-, 8-, 16-, and 32-states. The T2 trellis optimized for TC-CIOD improves the performance of differential TC-CIOD by 0.4 dB. For the sake of comparison, the CER performances of the nondifferential TC-STBC and TC-CIOD systems are also shown in Figure 3. As expected, the CER performances of nondifferential schemes have approximately 3 dB coding gain advantage compared to their differential counterparts. In Figure 4, the CER performances of the optimum differential TC-CIOD with trellis codes given in Table 1 are compared with those of 8-, 16-, and 32-state differential TC-STBC with optimum trellis codes proposed in [3, Table 1]. The perfectly interleaved multipath channel, 256 OFDM subcarriers, and perfect knowledge of the scaling coefficients $S_k^i$s, were assumed during the simulations. As seen from Figure 4, the proposed scheme considerably outperforms the differential two-symbol interleaved TC-STBC given in [3]. Using TC-CIOD instead of TC-STBC with aforementioned 8-, 16-, and 32-state trellis codes provides approximately 9.5 dB, 4 dB, and 3.5 dB SNR gain at the CER of $10^{-3}$.

Figure 5 shows the simulation results of the proposed differential TC-CIOD and reference two-symbol interleaved differential TC-STBC [3] with the same bandwidth efficiency over the COST 207 12-ray typical urban (TU) channel model [12]. The TC-CIOD and TC-STBC employ 4-state 8-PSK $R = 2/3$ trellis codes from Table 1 and [3], respectively. $K = 256$ OFDM subcarriers and OFDM symbol duration

| States | $h^0$ | $h^1$ | $h^3$ | $G_d$ | $G_t$ |
|--------|-------|-------|-------|-------|-------|
| 4      | 7     | 2     | 6     | 6     | 0.53  |
| 8      | 13    | 6     | 4     | 8     | 0.50  |
| 16     | 23    | 6     | 10    | 10    | 0.45  |
| 32     | 65    | 4     | 12    | 12    | 0.33  |

Figure 2: Codeword error probability upper bound of best trellis codes found for different values of $\theta$ ($R = 2/3$, 8-PSK, $E_b/N_0$ = 17 dB, $\kappa = 3$).
Figure 3: CER performances of TC-STBC and TC-CIOD OFDM schemes in a very slow varying fading channel ($K = 48, n_T = 2, n_R = 1$).

Figure 4: CER performances of differential TC-STBC and TC-CIOD OFDM schemes with 8-, 16-, and 32-state trellises in a very slow varying fading channel ($K = 256, n_T = 2, n_R = 1$).

Figure 5: CER performances of differential TC-STBC and TC-CIOD OFDM schemes with 4-state 8-PSK $R = 2/3$ trellis codes in COST 207 12-ray TU channel model ($K = 256, T_s = 128 \mu s, n_T = 2, n_R = 1, 2$ bps/Hz).

$T_s = 128 \mu s$ were selected during simulations. The CER performances with perfect knowledge (PK) of the scaling coefficients $S_k^f$ were simulated for normalized Doppler frequencies $f_{D,n} = 0.001$ and $f_{D,n} = 0.01$, that for OFDM symbol period $T_s = 128 \mu s$ and carrier frequency $f_c = 900$ MHz correspond to mobile terminal speeds $v = 9.37$ km/h and $v = 93.69$ km/h, respectively. Figure 5 shows that the high mobile terminal speeds cause an error floor due to the rapid change of channel weights. The simulations performed by estimating the scaling coefficients $S_k^f$ at the receiver by using (26) and (27) are indicated by the subchannel power estimation length $M$ in Figure 5. $M = 10$ and $M = 4$ were found to be optimum by exhaustive computer simulations for $f_{D,n} = 0.001$ and $f_{D,n} = 0.01$, respectively, under the considered channel conditions. When perfect channel interleaving is not considered, the selection of the channel interleaver $\alpha$ considerably affects the CER performances of TC-CIOD and TC-STBC systems. We performed the simulations for all possible block-type channel interleavers $\alpha$ and found that the performance of both systems improves when $2 \times 128$ block type channel interleaver is employed. Hence, all of the results given in Figure 5 are for $2 \times 128$ block channel interleaver.

Figure 5 shows that the perfect knowledge (PK) of the scaling coefficients $S_k^f$ provides approximately 2 dB and 4 dB SNR gain at the CER of $10^{-2}$ when $f_{D,n} = 0.001$ ($M = 10$) and $f_{D,n} = 0.01$ ($M = 4$), respectively. Note that we also simulated the TC-CIOD performance when scaling coefficients $S_k^f$ are estimated by using the previous decoder output in (13) to find $(|d_{jk}^f|^2 + |d_{jk+1}^f|^2)$ and used in (24). However, this method does not provide useful results due to error propagation. Figure 5 also shows that the proposed TC-CIOD scheme outperforms the reference TC-STBC [3] scheme by 4 dB at the CER of $10^{-2}$ and by 6 dB at the CER of $10^{-3}$ when $f_{D,n} = 0.001$. Additionally, the proposed scheme has a much lower error floor when channel weights are rapidly changing ($f_{D,n} = 0.01$).

Figure 6 shows the CER performances of the proposed differential TC-CIOD and the reference two-symbol interleaved differential TC-STBC [3] with 8-state 8-PSK $R = 2/3$ trellis codes from Table 1 and [3], respectively. The $2 \times 128$ block-type channel interleaver $\alpha$ is employed in all systems.
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