Exploiting Trust Degree for Multiple-Antenna User Cooperation

Mingxiong Zhao, Jong Yeol Ryu, Member, IEEE, Jemin Lee, Member, IEEE, Tony Q. S. Quek, Senior Member, IEEE, and Suili Feng, Member, IEEE

Abstract—For a user cooperation system with multiple antennas, we consider a trust degree based cooperation techniques to explore the influence of the trustworthiness between users on the communication systems. For the system with two communication pairs, when one communication pair achieves its quality of service (QoS) requirement, they can help the transmission of the other communication pair according to the trust degree, which quantifies the trustworthiness between users in the cooperation. For given trust degree, we investigate the user cooperation strategies, which include the power allocation and precoder design that maximize the expected achievable rates while guaranteeing the QoS requirement. For a SIMO case, we resort to semidefinite relaxation (SDR) technique and block coordinate update (BCU) method to solve the corresponding problem, and guarantee the rank-one solutions at each step. For a MIMO case, as MIMO is the generalization of MISO and SIMO, the similarities among their problem structures inspire us to combine the methods from MISO and SIMO together to efficiently tackle MIMO case. Simulation results show that the trust degree information has a great effect on the performance of the user cooperation in terms of the expected achievable rate, and the proposed user cooperation strategies achieve high achievable rates for given trust degree.

Index Terms—Trust degree, cooperative transmission, beamforming, power allocation

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M. Zhao is with the School of Software, Yunnan University, Kunming, China and also with Key Laboratory in Software Engineering of Yunnan Province. Email: jimmyzmx@gmail.com, mx_zhao@ynu.edu.cn.

J. Y. Ryu is with the Department of Information and Communication Engineering, Gyeongsang National University, Tongyeong, Korea. Email: ryuuyj18@gmail.com.

J. Lee is with the Department of Information and Communication Engineering, Daegu Gyeongbuk Institute of Science and Technology (DGIST), Daegu, Korea. Email: jmllee@dgist.ac.kr.

T. Q. S. Quek is with Singapore University of Technology and Design, Singapore 487372 and also with the Department of Electronics Engineering, Kyung Hee University, Yongin-si, Gyeonggi-do, 17104, Korea. E-mail: tonyquek@sutd.edu.sg.

S. Feng is with the School of Electronics and Information Engineering, South China University of Technology, Guangzhou, China. Email: fengsl@scut.edu.cn.

I. INTRODUCTION

The cooperative communications have been introduced to improve the communication reliability and spectral efficiency to satisfy the growing demand for higher data rates in wireless networks. Generally, there are two ways of realizing cooperation: 1) utilizing fixed relay terminals to assist the communication between dedicated sources and their corresponding destinations, 2) allowing the mobile users to help each other as a relay for reliable communication [1]. For utilizing fixed relays, the fixed relays need to be installed in the network, which requires high infrastructure, operation and maintenance costs for the operators [2]. On the other hand, in the user cooperation, relays are the mobile users who have good channel conditions and low traffic demands, and they can help the communications without increasing the cost to the mobile operators [3]. Furthermore, many mobile relays exist in the network, so each relay needs to assist few users only and the average power used by each mobile relay for transmission of signals is much smaller than that of the fixed relays [2]. Due to these advantages, in the cooperative communications, the user cooperation techniques have been intensively investigated [2–7].

In the user cooperation, the proximity between users enables their direct communications via device-to-device (D2D) communication. For various communication networks, the techniques for the cooperative D2D communications have been proposed [4–7]. For a cellular network, where the cellular and D2D users coexist, the relaying scheme of D2D user to assist the downlink transmission of cellular user was proposed in [4]. The cooperative D2D communication between the femto and macro users of a heterogeneous network (HetNet) was proposed in [5], where the femto user overheard and forwarded the composite of desired and interference signals to improve the signal-to-interference ratio (SINR) of the macro user. In [6], it is shown that mobile terminals can save the energy by exploiting the good channel quality of short range cooperation for WiFi and WiMedia. In the existing literatures [4–7], the user cooperation techniques are designed based on their traffic demands and corresponding qualities of physical channels.

However, different from the cooperation using the fixed relays, in the user cooperation, the relationship between users, e.g., trust degree, may affect the performance of user cooperation [8]. The social relationship between users can be an important motivation for participating in cooperative communications. Users would be willing to help each other by...
consuming their own resources if they have close relationship in the social domain. Otherwise, it is not sure whether users will cooperate even though they have the good channels and low traffic demands. Furthermore, some users, generally not with close relationship, may discard the data of the other user during the cooperation due to either the selfish behavior to save its own resource or the malicious purpose to disconnect the communications. Therefore, in the user cooperation, the relationship among users should be taken into account as a key design parameter for the efficient cooperative communications.

Recently, the social relationship has been actively considered in the development of communication strategies such as \([9]–[10]\). With the consideration of social relationships and physical coupling among users, a social group utility maximization framework was developed to maximize the social group utility, which is a sum of individual utilities weighted by its social ties with other users in \([9]\) and \([10]\). The social relationship of nodes has also been considered to enhance the performance of D2D communication \([11]–[13]\). Specifically, social trust and social reciprocity, which were achieved by exchanging the altruistic actions among nodes, were utilized in the D2D relay selection \([11]\), while the social-aware D2D communication architecture was proposed by exploiting social networking characteristics for system design \([12]\). In \([13]\), a traffic offloading for D2D communications was optimized for given online and offline social relations.

The trustworthiness between nodes has also been exploited for the design of efficient cooperation strategies \([14]–[16]\). In these works, the trust degree, which quantifies a degree of trustworthiness between nodes, was used as one of key design parameters to develop cooperative relay frameworks for the single-input-single-output (SISO) \([14]\) and multiple-input-single-output (MISO) \([15]\) systems. In terms of communication confidentiality, it was also shown that the expected secrecy rate can be increased by exploiting the trust degree of untrustworthy node in \([16]\). The previous works showed that the trust information can improve the performance of the conventional systems, which are designed based on the physical parameters only.

In this paper, motivated by the strong interests in trust degree, we investigate the user cooperation techniques based on trust degree between two pairs of communication users with multiple antennas, i.e., Tu1-Ru1 and Tu2-Ru2. Different from the existing works that consider a simple system model, we consider the cooperation techniques with multiple antennas, including power allocation and precoder design. The transmit user, Tu2, helps Tu1 by forwarding the information of Tu1 when Tu2 has good channel quality to Ru2, enough to guaranteeing its own quality of service (QoS) requirement. The willingness of Tu2 in helping Tu1 is characterized by the trust degree, i.e., Tu2 helps with high probability when trust degree is high. To maximize the expected achievable rate at Ru1, we jointly design the transmission strategies at Tu1 and Tu2 for four different antenna configurations: 1) SISO case where all users equip a single antenna as a special case, 2) MISO case where only Tu1 equips multiple antennas, 3) SIMO case where only Tu2 equips multiple antennas, and 4) MIMO case where both Tu1 and Tu2 equip multiple antennas. For SISO case, we first present an optimal power allocation strategy at Tu2, which maximizes the expected achievable rate at Ru1 while guaranteeing QoS requirement at Ru2. For MISO case, we provide an optimal structure of beamformer at Tu1 as a linear combination of the weighted channel vectors. Then, based on the structure, we obtain the beamformer that maximizes an approximated expected achievable rate as a function of the trust degree and corresponding power allocation at Tu2. For SIMO case, to jointly optimize the beamformers of Tu2, we utilize semidefinite relaxation (SDR) technique and block coordinate update (BCU) method to solve the considered problem, and guarantee the rank-one solutions at each step. Furthermore, for MIMO case, the similarities among the problem structures of MISO, SIMO and MIMO cases inspire us to combine the design of beamformer at Tu1 from MISO and the alternative algorithm from SIMO together to jointly optimize the beamformers at Tu1 and Tu2 to maximize the expected achievable rate at Ru1.

The rest of paper is organized as follows. We first describe the trust degree and system model in Section \(\text{II}\). In Section \(\text{III}\) based on the trust degree, the optimal transmission strategies and beamforming design in terms of the expected achievable rate are derived for four different cases. Numerical results are presented in Section \(\text{IV}\) and conclusions are drawn in Section \(\text{V}\).

**Notation:** In this paper, lowercase and uppercase boldface letters represent vectors and matrices, respectively. The complex conjugate of \(x\) is denoted by \(\bar{x}\), the hermitian transpose and the trace of \(X\) are denoted by \(X^\dagger\) and \(\text{tr}(X)\). \(\Pi_X \triangleq X(X^\dagger X)^{-1} X^\dagger\) represents the orthogonal projection onto the column space of \(X\), and \(\Pi_X^\perp \triangleq 1 - \Pi_X\) denotes the orthogonal projection onto the orthogonal complement of the column space of \(X\). \(X \sim \mathcal{CN}(A, B)\) denotes the elements of \(X\) that follow independent complex Gaussian distribution with mean \(A\) and covariance \(B\).

## II. System Model

We consider a system with two communication pairs: Tu1-Ru1 and Tu2-Ru2, where Tu1 and Tu2 are equipped with \(N_1\) and \(N_2\) antennas, respectively, and the receivers have a single antenna. In our system model, if one node achieves its own QoS requirement, it can help the transmission of the other node by using residual resource. Here, we assume that the node decides whether to help the transmission or not based on the trust degree, which measures the trustworthiness between nodes. Without loss of generality, Tu2 helps the transmission of Tu1 according to the trust degree between two user pairs. In the following subsections, we first briefly give an introduction of trust degree, and describe the system model in details.

### A. Trust Degree

With the explosive growth of online social networks such as WeChat and Facebook, a growing number of people are getting involved in online social interactions, and thus, the social relationship has been studied as an important parameter to investigate how the degree of closeness of social relationship between users affects their communication strategies \([9]–[11]\).
In the communication networks, the trust degree has been defined as a belief level that one node can put on another node for a certain action according to previous direct or indirect information, obtained from observations of behavior [17], [18]. Hence, in the cooperative communication systems, the trust degree can be interpreted as the degree that reveals how much a node is willing to help the communication of the other node [14], [15]. Similarly, in our system model, the trust degree between Tu1 and Tu2, α, is defined by the probability that Tu2 helps the transmission of Tu1 and thus, α is a value in range of 0 ≤ α ≤ 1.

In the previous literatures, the trust degree has been evaluated and quantified by various ways [8], [17]–[19]. The trust degree can be evaluated by the observations of the previous behaviors of the node [17]–[21]. In [20], [21], the trust degree is determined using Bayesian framework. In the Bayesian framework, the trust degree is given by the ratio of the observations of the positive behavior among total observations, where the positive behavior is that the node behaves in the predefined way of the network. Similar to [21], in our cooperative communication systems, the positive behavior is defined by that Tu2 helps the transmission of Tu1 and hence, Tu1 can estimate the trust degree based on the historical observations of the positive behavior of Tu2. The trust degree can also be updated according to new observations. However, when the number of observations is sufficiently large, the trust degree will have ignorable change according to new observation and it will be more like a constant. Therefore, in our system model, we assume that the trust degree is unchanged during the transmission.

In the user cooperative communications, the user may not help the other user’s transmission due to either the selfish behavior to save its own resource or the malicious purpose to disconnect the communication of that user. For the case of the malicious purpose, the malicious user lets the other user know that he will help the transmission. However, the malicious user can intercept or drop the data from the other user. Therefore, for the case that the users are not trustworthy, each user designs the transmission strategy based on the trust degree.

**B. System Description**

In this paper, we consider the cooperative communication system, where Tu2 can help the transmission of Tu1; if its corresponding receiver Ru2 achieves its QoS, as shown in Fig. 1. In our system, Tu2 decides whether to help the transmission of Tu1 based on the trust degree, α, which is defined as the probability that Tu2 cooperates with Tu1 as a relay node. Hence, for given α we design the optimal transmission strategy at Tu1 and the power allocation for cooperation at Tu2 to maximize an expected achievable rate while guaranteeing the QoS requirement at Ru2. Once Tu2 decides to help the transmission of Tu1, Tu2 determines the portion of transmission power β for relaying information and the portion for its own data transmission. The channels from Tu1 to Tu2, Ru1, and Ru2 are defined by $H_0 \in \mathbb{C}^{N_2 \times N_1}$, $h_1 \in \mathbb{C}^{N_1 \times 1}$, and $h_{12} \in \mathbb{C}^{N_1 \times 1}$, respectively, and they follow a complex Gaussian distribution with zero mean and covariances, $\sigma^2_{h_0}I_{N_2 \times N_1}$, $\sigma^2_{h_1}I_{N_1}$, and $\sigma^2_{h_{12}}I_{N_1}$, respectively. The channels from Tu2 to Ru1 and Ru2 are also defined by $h_{21} \in \mathbb{C}^{N_2 \times 1}$ and $h_2 \in \mathbb{C}^{N_2 \times 1}$, which have the covariances $\sigma^2_{h_{21}}I_{N_2}$ and $\sigma^2_{h_2}I_{N_2}$, respectively. Notice that when Tu2 helps the transmission of Tu1, Ru1 can estimate the channel from Tu2, $h_{21}$, and then, Ru1 reports the channel estimation of $h_{21}$ to Tu1 by feedback channel.

![Cooperative communication system](image)

**Fig. 1.** Cooperative communication system

In our system model, the data transmission operates in time-division mode, where Tu1 and Tu2 transmit their own data at $t_1$ and $t_2$, respectively, and $t_1$ and $t_2$ are assigned to be orthogonal with $t_1 = t_2 = t$. In the time slot $t_1$, Tu1 transmits the information-carrying symbol $x_1$ with $E[x_1^2] = 1$ to Ru1 and during $t_1$, Ru2 and Tu2 can also listen to $x_1$. To efficiently transmit data, Tu1 designs the transmit beamformer, $w_1 \in \mathbb{C}^{N_1 \times 1}$, which satisfies $w_1^H w_1 = P_1$ and $P_1$ is the maximum transmit budget at Tu1, and uses it for transmission in $t_1$. Hence, the received signals at Ru1, Ru2, and Tu2 in the time slot $t_1$ are respectively given by

$$y_{Ru1,1} = h_1^H w_1 x_1 + n_{Ru1,1},$$
$$y_{Ru2,1} = h_{12}^H w_1 x_1 + n_{Ru2,1},$$
$$y_{Tu2} = H_0^T w_1 x_1 + n_{Tu2},$$

where $n_{Ru1,1}$ and $n_{Ru2,1}$ represent the additive white Gaussian noise (AWGN) associated with Ru1 and Ru2, which follow the complex Gaussian distribution with zero mean and variance $\sigma^2$ (i.e., $\mathcal{CN}(0, \sigma^2)$), and $n_{Tu2} \in \mathbb{C}^{N_2 \times 1}$ is the AWGN at Tu2 with $\mathcal{CN}(0, \sigma^2 I_{N_2})$. We assume that if Ru2 can decode the data from the received signal, $y_{Ru2,1}$, in the time slot $t_1$, Ru2 can use it as side information for further improving the performance by interference cancellation in $t_2$.

In the time slot $t_2$, Tu2 transmits $x_2$ with $E[x_2^2] = 1$ to Ru2. Here, Tu2 has its traffic demand $Q$ and if Ru2 achieves $Q$ by the part of total power budget, Tu2 can help the transmission of Tu1 using residual power. According to the trust degree, Tu2 relays the data of Tu1 to Ru1 via decode-and-forward (DF) based relaying with probability $\alpha$ and otherwise, Tu2 transmits its own data only. Thus, the transmitted signal from Tu2 can be represented by

$$x_{Tu2} = w_{21} x_1 + w_{22} x_2,$$
where \( w_{21} \) and \( w_{22} \) are the transmit beamformers at Tu₂ for \( x_1 \) and \( x_2 \), respectively, and they are designed to satisfy the power constraint as

\[
w_{21}^\dagger w_{21} + w_{22}^\dagger w_{22} \leq P_2, \tag{5}\]

where \( P_2 \) is the maximum transmit budget at Tu₂ and \( w_{21}^\dagger w_{21} = \beta P_2 \). If Tu₂ does not help Tu₁ with probability \( 1 - \alpha \), Tu₂ does not allocate the power for \( x_1 \) such as \( \beta = 0 \). Therefore, the received signals at Ru₁ and Ru₂ in the time slot \( t_2 \) are respectively given by

\[
y_{R_{u_2},2} = h_{22}^\dagger x_{T_{u_2}} + n_{R_{u_2},2} = h_{22}^\dagger w_{22} x_2 + n_{R_{u_2},2}, \tag{6}\]

\[
y_{R_{u_1},2} = h_{21}^\dagger x_{T_{u_2}} + n_{R_{u_1},2} = h_{21}^\dagger w_{21} x_1 + n_{R_{u_1},2}, \tag{7}\]

where \( n_{R_{u_1},2} \) and \( n_{R_{u_2},2} \) are the AWGN at Ru₁ and Ru₂ with \( CN(0, \sigma^2) \) in \( t_2 \), respectively. If Ru₂ successfully decodes the data from Tu₁ in the time slot \( t_1 \), Ru₂ can subtract it from the received signal in \( t_2 \) by applying the successive interference cancellation (SIC). For this case, after applying SIC, the received signal at Ru₂ in \( t_2 \) can be rewritten as

\[
y_{R_{u_2},2} = h_{21}^\dagger w_{22} x_2 + n_{R_{u_2},2}. \tag{8}\]

Otherwise, Ru₂ has to decode \( x_2 \) by treating the signal related to \( x_1 \) as the noise.

In this paper, we only consider the case that Ru₂ can always achieve its QoS requirement, \( Q \), for given power budget, \( P_2 \), and hence, the QoS requirement at Ru₂ is given in the range of \( 0 \leq Q \leq Q_{\text{max}} \), where the maximum QoS requirement is \( Q_{\text{max}} = \frac{1}{2} \log (1 + \frac{\| h_{21}^\dagger \|^2}{\| h_{22}^\dagger \|^2} ) \), which is achieved by the maximal-ratio transmission (MRT) at Tu₂, \( w_{2}^{\text{mrt}} = \sqrt{\frac{P_2}{\| h_{22}^\dagger \|^2}} \) with maximum power \( P_2 \), and \( \frac{1}{2} \) is from the fact that the transmission takes place in two time slots.

### III. USER COOPERATION BASED ON TRUST DEGREE

In this section, for given trust degree, we provide the cooperation strategy, which includes the transmission beamformer at Tu₁ and the power allocation for cooperation at Tu₂, to maximize the expected achievable rate at Ru₁ while guaranteeing QoS requirement at Ru₂. We derive the optimal transmission strategies for three cases: 1) SISO case \( (N_1 = N_2 = 1) \), 2) MISO case \( (N_1 \geq 2, N_2 = 1) \), 3) SIMO case \( (N_1 = 1, N_2 \geq 2) \), and 4) MIMO case \( (N_1 \geq 2, N_2 \geq 2) \). We first define the event of cooperation as \( \mathcal{E} \), where \( \mathcal{E} = 1 \) and \( \mathcal{E} = 0 \) stand for the events that Tu₂ helps and does not help the transmission, respectively. Thus, \( \mathcal{E} \) is a Bernoulli random variable with \( P(\mathcal{E} = 1) = \alpha \) and \( P(\mathcal{E} = 0) = 1 - \alpha \).

As a performance metric, for given trust degree \( \alpha \), we use the expected achievable rate with respect to the possible cooperation events, defined as

\[
R_{\text{Ru}_1} = E \{ \tilde{R}_{\text{Ru}_1} \}, \tag{8}\]

where \( \tilde{R}_{\text{Ru}_1} \) is an achievable rate at Ru₁.

#### A. SISO case \( (N_1 = N_2 = 1) \)

We first consider a simple SISO case that Tu₁ and Tu₂ have a single antenna \( (N_1 = N_2 = 1) \). We define the gains of all channels as

\[
g_0 = |h_0|^2, g_1 = |h_1|^2, g_2 = |h_2|^2, g_{12} = |h_{12}|^2, g_{21} = |h_{21}|^2. \tag{9}\]

When the channel condition between Tu₁ and Tu₂ is worse than the direct channel from Tu₁ to Ru₁ (i.e. \( g_0 \leq g_1 \)), the cooperation of Tu₂ cannot improve the achievable rate at Ru₁ due to DF relaying constraint [22]. Thus, in this case, the achievable rate at Ru₁ is achieved by the direct transmission from Tu₁. Therefore, since Tu₂ helps the transmission of Tu₁ with probability \( \alpha \), the expected achievable rate at Ru₁ is given by

\[
R_{\text{Ru}_1}(\beta) = \begin{cases} 
\tilde{R}_{\text{Ru}_1}(\beta), & \text{if } g_0 > g_1, \\
\frac{1}{2} \log_2 (1 + \rho_1 g_1), & \text{otherwise}, 
\end{cases} \tag{10}\]

where \( \tilde{R}_{\text{Ru}_1}(\beta) \) is given by

\[
\tilde{R}_{\text{Ru}_1}(\beta) = \frac{\alpha}{2} \min \left\{ \log_2 \left( 1 + \rho_1 g_1 + \frac{\beta \rho_2 g_21}{(1 - \beta) \rho_2 g_{21} + 1} \right), \log_2 (1 + \rho_1 g_0) \right\} + \frac{1 - \alpha}{2} \log_2 (1 + \rho_1 g_1), \tag{11}\]

where \( \rho_1 = \frac{g_21}{g_0} \) and \( \rho_2 = \frac{g_21}{g_1} \). From the observations of (10) and (11) in SISO case, we deduce that \( R_{\text{Ru}_1} \) is an increasing function of \( \rho_1 \). Therefore, as \( \rho_1 \) grows, the value of \( R_{\text{Ru}_1} \) is increasing, and the maximum transmit power \( P_1 \) is always optimal.

The first term of (11) denotes the achievable rate at Ru₁ when Tu₂ helps transmission of Tu₁ with probability \( \alpha \), and it is bounded by the minimum of achievable rates at Tu₂ and Ru₁ due to the constraint of DF relaying [22]. The second term of (11) represents the achievable rate achieved by direct transmission from Tu₁ to Ru₁ when Tu₂ does not help transmission of Tu₁ with probability \( 1 - \alpha \). In the SISO case, we can see that \( \beta \) can be determined independently from \( \alpha \). The power allocation \( \beta \) has to be jointly determined with the transmit strategy at Tu₁, which is related to \( \alpha \). However, in SISO case, Tu₂ does not design the transmit beamformer and hence, \( \beta \) can be determined independently with \( \alpha \) to maximize the rate achieved when Tu₂ cooperates with Tu₁.

We define the part related with \( \beta \) in (11) as \( Q_{\text{Ru}_1}(\beta) \), given by

\[
Q_{\text{Ru}_1}(\beta) = \frac{1}{2} \min \left\{ \log_2 \left( 1 + \rho_1 g_1 + \frac{\beta \rho_2 g_{21}}{(1 - \beta) \rho_2 g_{21} + 1} \right), \log_2 (1 + \rho_1 g_0) \right\}. \tag{12}\]

According to \( \beta \), \( Q_{\text{Ru}_1}(\beta) \) can be rewritten by

\[
Q_{\text{Ru}_1}(\beta) = \begin{cases} 
\frac{1}{2} \log_2 (1 + \rho_1 g_0), & \text{if } \beta_0 \leq \beta \leq 1, \\
\frac{1}{2} \log_2 \left( 1 + \rho_1 g_1 + \frac{\beta \rho_2 g_{21}}{(1 - \beta) \rho_2 g_{21} + 1} \right), & \text{otherwise}, 
\end{cases} \tag{13}\]

where \( \beta_0 = 1 - \frac{\rho_2 g_{21} - \rho_1 (g_0 - g_1)}{\rho_2 g_{21} (1 + \rho_1 g_1)} \).

For the case that Tu₂ helps the transmission of Tu₁, if the receiving rate at Ru₂ in \( t_1 \), \( R_{12} = \frac{1}{2} \log_2 (1 + \rho_1 g_{12}) \geq Q_{\text{Ru}_1}(\beta) \), Ru₂ can decode \( x_1 \) in \( t_1 \) and hence, Ru₂ can apply SIC to eliminate the effect of \( x_1 \). Since the QoS requirement, \( Q \), is always satisfied when Tu₂ does not help the transmission of Tu₁, we only consider the case that Tu₂ helps the transmission of Tu₁ to design the transmission strategy.
in its received signal in $t_2$. Therefore, the achievable rate at $R_u_2$ is given by

$$R_{R_u_2}(\beta) = \begin{cases} \frac{1}{2} \log_2 \left( 1 + (1 - \beta_0) \rho_2 g_2 \right), & \text{if } R_1 \geq Q_{T_u_1}(\beta), g_0 > g_1, \\ \frac{1}{2} \log_2 \left( 1 + (1 - \beta_1) \rho_2 g_2 \right), & \text{if } R_1 < Q_{T_u_1}(\beta), g_0 > g_1, \end{cases}$$

(14)

where $R_{R_u_2}(\beta)$ and $R_{R_u_2}^{\text{NSC}}(\beta)$ are given, respectively, by

$$R_{R_u_2}(\beta) = \frac{1}{2} \log_2 \left( 1 + (1 - \beta) \rho_2 g_2 \right),$$

$$R_{R_u_2}^{\text{NSC}}(\beta) = \frac{1}{2} \log_2 \left( 1 + (1 - \beta_0) \rho_2 g_2 \right).$$

(15)

For the SISO case, the optimal power allocation of $T_u_2$ that maximizes the expected achievable rate at $R_u_1$ while guaranteeing QoS requirement at $R_u_2$ is obtained by the following problem

$$\textbf{P1:} \quad \max_{0 \leq \beta \leq 1} R_{R_u_2}(\beta)$$

subject to $R_{R_u_2}(\beta) \geq Q$,

(16a)

(16b)

where $R_{R_u_2}(\beta)$ is given in (10) and $Q$ is the QoS requirement at $R_u_2$.

For the MISO case, the optimal power allocation of $T_u_2$ for cooperative transmission in the following theorem.

**Theorem 1:** For given channels and QoS requirements, $Q$, the optimal power allocation of $T_u_2$ for cooperative transmission that maximizes the expected achievable rate at $R_u_1$ is obtained by

$$\beta^* = \begin{cases} \beta_1 Q_1, & \text{if } g_1 \geq Q, r_1, \\ \beta_2, & \text{if } g_0 > g_1 \geq r_1, \\ \beta_2, & \text{otherwise}, \end{cases}$$

(17)

where $\beta_1 Q_1$, $\beta_2$, and $\beta_1$ are given, respectively, by

$$\beta_1 Q_1 = 1 - \frac{4Q - 1}{\rho_2 g_2},$$

$$\beta_2 = \left( 1 - \frac{4Q - 1}{\rho_2 g_2} \right)^{-1} - Q,$$

$$\beta_1 = 1 - \frac{\rho_2 g_2 (1 + \rho_1 (g_1 - g_1))}{\rho_2 g_2 (1 + \rho_1 (g_1 - g_1))},$$

(18)

(19)

(20)

and $r_1$, $r_2$, and $r_3$ are given by

$$r_1 = \left[ \frac{1}{2} \log_2 \left( 1 + (1 - \beta_0) \rho_2 g_2 \right) \right]^+, \quad r_2 = \left[ \frac{1}{2} \log_2 \left( 1 + (1 - \beta_1) \rho_2 g_2 \right) \right]^+, \quad r_3 = \left[ \frac{1}{2} \log_2 \left( \frac{1}{1 + \rho_2 g_2} \right) \right]^+.$$ 

(21)

**Proof:** The proof is presented in Appendix A.

**Remark 1:** From Theorem 1 we observe that the optimal $\beta$ that maximizes the expected achievable rate at $R_u_1$ while guaranteeing QoS of $R_u_2$ is mainly determined by channel quality from $T_u_1$ to $R_u_2$, $g_{12}$ and QoS of $R_u_2$, $Q$. Since $T_u_2$ can assign more power for cooperation when $R_u_2$ can apply SIC than that when $R_u_2$ cannot apply SIC, the optimal power allocation is mainly determined by the parameters that decide whether SIC is applicable at $R_u_2$, $g_{12}$ and $Q$.

For the case that the channel quality from $T_u_1$ to $R_u_2$ is good such as $g_{12} \geq g_0$ and $g_{12} \geq g_1$, since we have $r_1 \geq r_2$ from $\beta_0 \leq \beta_1$, the optimal power allocation for cooperation is given by $\beta^* = \beta_1$, for all $Q$. In this case, for all $Q$, $R_u_2$ can apply SIC due to good channel quality between $T_u_1$ and $R_u_2$.

Contrarily, when the channel quality from $T_u_1$ to $R_u_2$ is poor such as $g_0 > g_{12}$ and $g_1 > g_{12}$, $R_u_2$ cannot apply SIC for all $Q$ and hence, the $T_u_2$ allocates the minimum power for cooperation as $\beta^* = \beta_2$. For the moderate quality of $g_{12}$ as $g_0 > g_{12} \geq g_1$, the power allocation is also determined according to QoS requirement, $Q$. In this case, since $R_u_2$ cannot apply SIC due to $g_1$ and $Q$, $\beta$ has to be controlled to apply SIC. First, for the case that QoS requirement is small such as $Q \leq r_3$, since $T_u_2$ can allocate large power for cooperation without SIC as $\beta^* = \beta_2$, which is a decreasing function with $Q$, SIC by reducing $\beta$ is not beneficial. For moderate QoS as $r_2 > Q > r_3$, $T_u_2$ controls $\beta$ as $\beta^* = \beta_1$, which is a constant for given channels, to apply SIC. In the moderate $Q$, the power for cooperation is not decreased even if $Q$ increases. However, when QoS requirement is high as $Q \geq r_2$, the power allocation for cooperation is decreased according to $Q$ as $\beta^* = \beta_2$ to guarantee high QoS by applying SIC.

**B. MISO case ($N_1 \geq 2$ and $N_2 = 1$)**

In this subsection, we consider the MISO case, where $T_u_1$ equips with multiple antennas (i.e., $N_1 \geq 2$) but $T_u_2$ has a single antenna (i.e., $N_2 = 1$). Hence, in the MISO case, we jointly design the transmit beamformer at $T_u_1$, $w_1$, and the power allocation for cooperation at $T_u_2$, $\beta$, to maximize the expected achievable rate at $R_u_1$ while guaranteeing the QoS requirement at $R_u_2$. For the MISO case, the expected achievable rate at $R_u_1$ is given by

$$R_{R_u_1}(w_1, \beta) = \begin{cases} \frac{1}{2} \log_2 \left( 1 + \rho_1 |h_2| w_1 \right)^2, & \text{if } \tilde{g}_0 > \tilde{g}_1, \\ \frac{1}{2} \log_2 \left( 1 + \rho_1 |h_2| w_1 \right)^2 + \frac{\beta \rho_2 g_2}{(1 - \beta) \rho_2 g_2 + 1}, & \text{otherwise}, \end{cases}$$

(22)

where $\tilde{g}_0 = \|h_0\|^2$, $\tilde{g}_1 = \|h_1\|^2$ and $\tilde{R}_{R_u_1}(w_1, \beta)$ is given by

$$\tilde{R}_{R_u_1}(w_1, \beta) = \frac{\alpha}{2} \min \left[ \log_2 \left( 1 + \rho_1 |h_2| w_1 \right)^2 + \frac{\beta \rho_2 g_2}{(1 - \beta) \rho_2 g_2 + 1}, \right] \left[ \frac{1 - \alpha}{2} \log_2 \left( 1 + \rho_1 |h_2| w_1 \right)^2 \right].$$

(23)

The part related with both $w_1$ and $\beta$ in (23) is defined by

$$Q_{R_u_1}(w_1, \beta) = \frac{1}{2} \min \left[ \log_2 \left( 1 + \rho_1 |h_2| w_1 \right)^2 + \frac{\beta \rho_2 g_2}{(1 - \beta) \rho_2 g_2 + 1}, \right] \left[ \frac{1 - \alpha}{2} \log_2 \left( 1 + \rho_1 |h_2| w_1 \right)^2 \right].$$

(24)

Similar to SISO case, the expected achievable rate at $R_u_1$ is an increasing function of $P_1$ and hence, the maximum transmit power $P_1$ is always optimal for MISO case.

For the case that $T_u_2$ helps the transmission of $T_u_1$, if the receiving rate at $R_u_2$ in $t_1$, $R_{R_u_2}(w_1)$, is greater than $Q_{R_u_2}(w_1, \beta)$, i.e., $R_{R_u_2}(w_1) = \frac{1}{2} \log_2 \left( 1 + \rho_1 |h_2| w_1 \right)^2 \geq Q_{R_u_2}(w_1, \beta)$, $R_u_2$ can decode $x_1$ in $t_1$ and hence, $R_u_2$ can
apply SIC to eliminate the effect of $x_1$ in its received signal in $t_2$. Therefore, the achievable rate at $R_{u_2}$ is given by

$$R_{u_2}(\beta) = \begin{cases} R_{NSIC}(\beta), & \text{if } R_{12}(w_1) \geq Q_{tu_1}(w_1, \beta), \bar{g}_0 > \bar{g}_1, \\ R_{NSIC}(\beta), & \text{if } R_{12}(w_1) < Q_{tu_1}(w_1, \beta), \bar{g}_0 > \bar{g}_1, \\ \frac{1}{2} \log_2(1 + p_2 g_2), & \text{otherwise}, \end{cases}$$

(25)

where $R_{NSIC}(\beta)$ and $R_{NSIC}(\beta)$ are given in [15].

For the MISO case, the optimal beamformer at $Tu_1$, $w_1$, and the power allocation for cooperation at $Tu_2$, $\beta$, that maximize the expected achievable rate at $R_{u_1}$ while guaranteeing QoS requirement at $R_{u_2}$ are obtained by solving the following joint optimization problem

$$\textbf{P2} : \max_{w_1, \beta} R_{u_1}(w_1, \beta) \quad \text{s.t. } R_{u_2}(\beta) \geq Q, \quad w_1 \perp 1.$$  

(26a)

We define the constant values for given channels, $v_1$, $v_2$ and $\phi_i$ as

$$v_1 \triangleq \|\Pi_{h_0} h_1\|^2, \quad v_2 \triangleq \|\Pi_{h_0} h_1\|^2, \quad \phi_i \triangleq \rho_1 \left(1 + \frac{1}{\rho_2 g_2}\right).$$

Then, the optimal structure of beamformer $w_1$ can be obtained by the following lemma.

**Lemma 1.** [15] Lemma 1] The optimal beamformer at $Tu_1$ that maximizes the expected achievable rate at $R_{u_1}$ can be represented by

$$w_1^{\text{opt}} = \sqrt{\eta} w_0 + \sqrt{1 - \eta} w_0^\perp,$$  

(27)

where $w_0 = \Pi_{h_0} h_1$, $w_0^\perp = \Pi_{h_0} h_1^\perp$, and $\eta$ is a constant in the range of $\frac{v_1}{v_1 + v_2} \leq \eta \leq 1$.

**Proof:** The proof of Lemma 1 can be referred to [15].

In Lemma 1 it is difficult to present $\eta$ that maximizes $R_{u_1}(w_1, \beta)$ in closed form and hence, the optimal $\eta$ should be found by exhaustive search. However, from the numerically obtained beamformer $w_1$, we cannot get the insight on the effect of trust degree on $R_{u_1}(w_1, \beta)$. Hence, in order to obtain $w_1$ in closed form, the approximated expected achievable rate at $R_{u_1}$ can be obtained by high signal-to-noise ratio (SNR) approximation [12] as

$$R_{u_1}(w_1, \beta) \approx \frac{\alpha}{2} \min \left\{ \log_2 \left( \frac{\rho_1 |h_1|^2 + \beta \rho_2 g_2}{1 - \beta \rho_2 g_2 + 1} \right) \right\} \left( \log_2 \left( \frac{\rho_1 |h_0|^2}{1 - \beta \rho_2 g_2 + 1} \right) \right) + \frac{1 - \alpha}{2} \log_2 \left( \frac{\rho_1 |h_1|^2}{1 - \beta \rho_2 g_2 + 1} \right).$$

(28)

For given $\beta$, the transmit beamformer that maximizes the approximated expected achievable rate at $R_{u_1}$ can be obtained by the following theorem.

**Theorem 2:** For given trust degree $\alpha$ and the power allocation $\beta$, the transmit beamformer of $Tu_1$ that maximizes $R_{u_1}(w_1, \beta)$ is given by

$$w_1^* = \sqrt{\eta^*} w_0 + \sqrt{1 - \eta^*} w_0^\perp,$$  

(29)

where $\eta^*$ is given by [30].

In order to obtain the closed-form beamformer, we adopt the high SNR approximation. However, it does not mean that we assume the high SNR configuration in our system model.

**Proof:** The proof is presented in Appendix B.

**Remark 2:** From Theorem 2 we can see that the direction of beamformer at $Tu_1$, $\eta$, is affected by both the trust degree, $\alpha$, and the power allocation for cooperation, $\beta$. First, when the power allocation for cooperation is very small such as $\beta < \beta_0$, the expected achievable rate enhancement from the cooperation with $Tu_2$ is small even if $\alpha$ is large. Hence, in this case, the beamformer at $Tu_1$ is designed to maximize the direct link from $Tu_1$ to $Ru_1$ regardless of $\alpha$. On the other hand, for high $\beta$ as $\beta > \beta_0$, since the expected achievable rate enhancement is large enough due to high $\beta$, the direction of beamformer mainly depends on the trust degree, $\alpha$. Hence, when $Tu_2$ helps the transmission of $Tu_1$ with high probability, equivalently $\alpha$ is high, the direction of beamformer is steered toward $h_0$ based on $\eta_2$, which is an increasing function with $\alpha$ to fully exploit the cooperation of $Tu_2$. Otherwise, the direction of beamformer has to be properly steered based both $\alpha$ and $\beta$. When $\alpha$ is relatively high compared to $\beta$ such as $\eta_2 > \eta_1(\beta)$, the direction of beamformer tends to be steered toward $h_1$ rather than $h_1$ and vice versa.

In Theorem 2, we show that beamformer $w_1$ can be represented according to $\beta$. Thus, the joint optimization problem of $w_1$ and $\beta$ can be simplified by the optimization problem of a single parameter $\beta$, as

$$\textbf{P2} - 1 : \max_{\beta} R_{u_1}(\beta) \quad \text{s.t. } R_{u_2}(\beta) \geq Q.$$  

(33a)

For the general case, it is hard to directly obtain the optimal $\beta^*$ from $\textbf{P2} - 1$ because $\textbf{P2} - 1$ is non-convex with respect to $\beta$. Thus, in the following corollaries, we obtain the optimal power allocation for cooperation for some special cases. In the following corollaries, we assume $\rho_1 = \rho_2$ for simplicity.

**Corollary 1:** When the channel from $Tu_1$ to $Tu_2$ is very strong such as $\bar{g}_0 \geq \bar{g}_0 = \frac{2(\bar{g}_1 - \bar{g}^2_2)}{\bar{g}_0}$, the optimal power allocation for cooperation that maximizes the approximated achievable rate at $R_{u_1}$ can be represented by

$$\beta^* = \begin{cases} \min \{ \tilde{\beta}_2, \tilde{\beta}_Q \}, & \text{if } v_3 \geq \tilde{g}_1^2, \\ \tilde{\beta}_Q, & \text{otherwise}, \end{cases}$$

(34)

where $\tilde{\beta}_2$ and $v_3$ are given by

$$\tilde{\beta}_2 = \frac{(v_3 - \tilde{g}_1^2) \phi_1}{\tilde{g}_1 + (v_3 - \tilde{g}_1^2) \rho_1}, \quad v_3 = |h_{12}^\perp h_1|^2,$$  

(35)

and $\tilde{\beta}_Q$ and $\tilde{\beta}_Q$ are given in [18] and [19], respectively.

**Proof:** For $\bar{g}_0 \geq \bar{g}_0$, since we have $\beta \geq \beta_0$, for all feasible $\beta$ in $0 \leq \beta \leq 1$, the optimal beamformer is given by $w_1^*(\eta_1) = \sqrt{\eta_1} w_0 + \sqrt{1 - \eta_1} w_0^\perp$, and the approximated achievable rate at $R_{u_1}$ is represented by [33] with $\eta_1$. For this case, the condition to apply SIC at $R_{u_2}$ is

$$R_{12}(w_1^*) = \frac{1}{2} \log_2 \left(1 + \frac{\rho_1 v_3}{\bar{g}_1}\right) \geq \frac{1}{2} \log_2 \left(1 + \frac{\rho_1 (\tilde{g}_1 + m_1(\beta))}{\tilde{g}_1}\right) = Q_{nu_2}(w_1^*, \beta) \Rightarrow \beta \leq \tilde{\beta}_2.$$  

(36)
requirement is obtained by \( \beta \leq \min \{ \tilde{\beta}_2, \beta_{Q_1} \} \). Since (B.4) is an increasing function of \( \beta \), the optimal power allocation is obtained by \( \beta^* = \beta_{Q_2} \).

Otherwise, Ru2 cannot apply SIC for all \( \beta \) and hence, the power allocation to guarantee QoS without SIC is obtained by \( \beta^* = \beta_{Q_2} \).

**Corollary 2:** When the channel from Tu1 to Tu2 is very weak such as \( \tilde{g}_0 < v_1 \), the optimal power allocation for cooperation that maximizes the approximated achievable rate at Ru1 can be represented by

\[
\beta^* = \begin{cases} 
\beta_{Q_1}, & \text{if } v_4 \geq v_5, \\
\beta_{Q_2}, & \text{otherwise},
\end{cases}
\]

where \( v_4 \) and \( v_5 \) are given by

\[
v_4 = |h_{12}^\dagger w_1^* (\eta_2)|^2, \quad v_5 = |h_0^\dagger w_1^*(\eta_2)|^2.
\]

**Proof:** If \( \tilde{g}_0 < v_1 \), for all feasible \( \beta \) in \( 0 \leq \beta \leq 1 \), the optimal beamformer is given by \( w_1^*(\eta_2) = \sqrt{\eta_2} w_0 + \sqrt{1-\eta_2} w_0 \). Hence, the remaining part of the proof of Corollary 2 can be obtained in the similar way to the proof of Corollary 1.

From Lemma 1, Theorem 2 and high SNR approximation, the joint optimization problem to design beamformer at Tu1, \( w_1 \), and the power allocation at Tu2, \( \beta \), is simplified into the optimization problem with a single parameter \( \beta \). For some special case, the optimal \( \beta \) can be obtained in a closed form. For general case, the optimal power allocation for cooperation can be obtained by one dimensional search from 0 to 1, which is much simpler than solving the joint optimization problem.

**C. SIMO case (\( N_1 = 1 \) and \( N_2 \geq 2 \))**

In this subsection, we consider the SIMO case, where Tu1 equips with a single antenna (i.e., \( N_1 = 1 \)) but Tu2 is equipped with \( N_2 \geq 2 \) antennas and hence, the SIMO channel is formed between Tu1 and Tu2. For the SIMO case, we jointly design the transmit beamformers at Tu2, \( w_21 \) and \( w_22 \), to maximize the expected achievable rate at Ru1. For this case, the expected achievable rate at Ru1 in SIMO case is given by

\[
R_{tu1}(w_21, w_22) = \begin{cases} 
\bar{R}_{tu1}(w_21, w_22), & \text{if } \tilde{g}_0 > g_1, \\
\frac{1}{2} \log_2 (1 + \rho_1 g_1), & \text{otherwise},
\end{cases}
\]

where \( \tilde{g}_0 > g_1 \) means the channel condition between Tu1 and Tu2 is better than the direct channel from Tu1 to Ru1, thus the cooperation of Tu2 can improve the achievable rate at Ru1, and \( \bar{R}_{tu1}(w_21, w_22) \) is given by

\[
\bar{R}_{tu1}(w_21, w_22) = \frac{\alpha}{2} \log_2 \left[ 1 + \rho_1 g_1 + \frac{|h_1^\dagger w_21|^2}{|h_1^\dagger w_22|^2 + \sigma^2} \right],
\]

\[
\log_2 (1 + \rho_1 \tilde{g}_0) + \frac{1-\alpha}{2} \log_2 (1 + \rho_1 g_1).
\]

We define the first term in (39) as

\[
Q_{tu1}(w_21, w_22) = \frac{1}{2} \min \left[ \log_2 \left( 1 + \rho_1 g_1 + \frac{|h_1^\dagger w_21|^2}{|h_1^\dagger w_22|^2 + \sigma^2} \right), \log_2 (1 + \rho_1 \tilde{g}_0) \right].
\]

For the case that Tu2 helps the transmission of Tu1, if the rate achieved at Ru2 in \( t_1 \), \( R_{tu2} \), is greater than \( Q_{tu1}(w_21, w_22) \), i.e., \( R_{tu2} = \frac{1}{2} \log_2 (1 + \rho_1 g_1) \geq Q_{tu1}(w_21, w_22) \), Ru2 can decode \( x_1 \) in \( t_1 \) and thus, Ru2 can employ SIC to eliminate the effect of \( x_1 \) in its received signal in \( t_2 \). Therefore, the achievable rate at Ru2 is given by

\[
R_{tu2}(w_21, w_22) = \begin{cases} 
R_{tu2}^{SIC}(w_21, w_22), & \text{if } R_{tu2} \geq Q_{tu1}(w_21, w_22), \\
\frac{1}{2} \log_2 (1 + \rho_2 \tilde{g}_2), & \text{otherwise},
\end{cases}
\]

where \( R_{tu2}^{SIC}(w_21, w_22) \) and \( R_{tu2}^{NSIC}(w_21, w_22) \) are respectively given by

\[
R_{tu2}^{SIC}(w_21, w_22) = \frac{1}{2} \log_2 \left( 1 + \frac{|h_{12}^\dagger w_21|^2}{\sigma^2} \right),
\]

\[
R_{tu2}^{NSIC}(w_21, w_22) = \frac{1}{2} \log_2 \left( 1 + \frac{|h_1^\dagger w_21|^2}{|h_1^\dagger w_22|^2 + \sigma^2} \right).
\]

For the SIMO case, to maximize the expected achievable rate at Ru1 while guaranteeing QoS of Ru2, the beamformers at Tu2, \( w_21 \) and \( w_22 \), are jointly optimized by the following problem

\[
P_3: \max_{w_21, w_22} R_{tu1}(w_21, w_22)
\]

\[
s.t. R_{tu2}(w_21, w_22) \geq Q,
\]

\[
w_21^\dagger w_21 + w_22^\dagger w_22 \leq P_2.
\]
where the power allocation at Tu2 is embedded in the beamformer design.

By considering the condition whether Ru2 applies SIC, P3 in (43) can be divided by two subproblems. For the case that Ru2 applies SIC, the optimization problem P3 – 1 is given by

\[ \text{P3 – 1 : max}_{w_1, w_2} \quad R_{tu_1}(w_1, w_2) \]
\[ \text{s.t.} \quad \frac{1}{2} \log_2 \left( 1 + \frac{|h_1 w_2|^2}{\sigma^2} \right) \geq Q, \]
\[ w_1^+ w_1 + w_2^+ w_2 \leq P_2, \]

where constraint (44b) is the condition that Ru2 decodes the data from Tu1 in t1 and applies SIC to cancel it in t2. Similarly, the case that Ru2 does not apply SIC, P3 – 2, is given by

\[ \text{P3 – 2 : max}_{w_1, w_2} \quad R_{tu_1}(w_1, w_2) \]
\[ \text{s.t.} \quad \frac{1}{2} \log_2 \left( 1 + \frac{|h_1 w_2|^2}{\sigma^2} \right) \geq Q, \]
\[ w_1^+ w_1 + w_2^+ w_2 \leq P_2. \]

For the relaying case, i.e., \( \tilde{g}_0 > g_1 \), from (39), \( R_{tu_1}(w_1, w_2) \) can be presented by two different forms according to \( w_1 \) and \( w_2 \) due to DF relaying constraint. Therefore, both P3 – 1 and P3 – 2 can be further divided into two subproblems with respect to \( R_{tu_1}(w_1, w_2) \). Hence, we can obtain the solution of P3 in (43) by choosing the best solution from the solutions of four subproblems. Since all subproblems can be solved in a similar way, here we focus on one of the four subproblems. In the following, we consider the subproblem P3 – 21, where Ru2 does not apply SIC in t2 and \( Q_{tu_1}(w_1, w_2) \) is determined by

\[ Q_{tu_1}(w_1, w_2) = \frac{1}{2} \log_2 \left( 1 + \rho_1 g_1 + \frac{|h_1 w_2|^2}{|h_1 w_2|^2 + \sigma^2} \right). \]

Therefore, for this case, the subproblem P3 – 21 is represented by

\[ \text{P3 – 21 : max}_{w_1, w_2} \quad \frac{|h_1 w_2|^2}{|h_2 w_2|^2 + \sigma^2} \]
\[ \text{s.t.} \quad \log_2 \left( 1 + \rho_1 g_1 + \frac{|h_1 w_2|^2}{|h_2 w_2|^2 + \sigma^2} \right) \geq Q, \]
\[ w_1^+ w_1 + w_2^+ w_2 \leq P_2, \]

where the constraint (47b) is the condition that Ru2 cannot apply SIC in t2 and the constraint (47c) is the condition to satisfy (45) referring to (40). Notice that constant \( \frac{1}{2} \) is ignored in (47c) without changing the property of the problem. From the observation of P3 – 21, the second term of the objective function, i.e., \( (1-\alpha) \log_2(1+\rho_1 g_1) \), is a constant, which can be ignored to obtain the solution. Then, we can see that for given \( \alpha \), the right hand sides (RHSs) of constraints (47a) and (47c) are equal to the objective function of (47a) to be maximized, i.e., \( \log_2 \left( 1 + \rho_1 g_1 + \frac{|h_1 w_2|^2}{|h_2 w_2|^2 + \sigma^2} \right) \). Hence, we note that if the problem P3 – 21 is feasible, the constraint (47b) is always hold. Otherwise, P3 – 21 is infeasible and the solution is obtained by the other subproblems. In addition, if the constraint (47c) is not hold, i.e., \( \log_2 \left( 1 + \rho_1 g_1 + \frac{|h_1 w_2|^2}{|h_2 w_2|^2 + \sigma^2} \right) \), we observe that the expected achievable rate is bounded by constant, which is independent from \( w_1 \) and \( w_2 \). In other word, Ru1 can achieve at least constant rate in (39) if the constraint (47c) is not hold. Hence, we can solve P3 – 21 without considering the constraint (47c) and after solving P3 – 21 without (47c), we can check whether the constraint (47c) is hold or not for the obtained solution. If (47c) is hold, the expected achievable rate at Ru1 is determined by the obtained solution and otherwise, the expected achievable rate is determined by the constant rate in (39). Therefore, by removing constraints (47b) and (47c), the subproblem P3 – 21 can be equivalently rewritten as

\[ \text{P3 – 21’ : max}_{w_1, w_2} \quad \frac{|h_1 w_2|^2}{|h_2 w_2|^2 + \sigma^2} \]
\[ \text{s.t.} \quad |h_2 w_2|^2 \geq (4^Q-1)|h_2 w_2|^2 + \sigma^2, \]
\[ w_1^+ w_1 + w_2^+ w_2 \leq P_2, \]

Since the problem P3 – 21’ is non-convex and \( w_1, w_2 \) are still coupled in the constraints, it is hard to directly obtain the solution of P3 – 21’ in its current form. To solve P3 – 21’ by decoupling \( w_1 \) and \( w_2 \), we apply block coordinate update (BCU) method to update \( w_1 \) (or \( w_2 \)) while fixing \( w_2 \) (or \( w_1 \)) at one iteration, and optimize \( w_2 \) (or \( w_1 \)) based on the newly updated \( w_1 \) (or \( w_2 \)) at the next iteration. Thus, the expected achievable rate at Ru1 is maximized by optimizing \( w_1 \) and \( w_2 \) iteratively. Employing the semidefinite relaxation (SDR) technique (23) and giving \( w_2 \) at k + 1-th iteration, the relaxation of (48c) can be rewritten as

\[ \text{P3 – 21’ – 1 : max}_{W_1 \geq 0} \quad \text{tr}(W_1 h_1 h_2^H) \]
\[ \text{s.t.} \quad |h_2 w_2|^2 \geq (4^Q-1)|h_2 w_2|^2 + \sigma^2, \]
\[ \text{tr}(W_1) + |w_2|^2 \leq P_2, \]

where we discard the constraint rank(\( W_1 \)) = 1. The relaxed problem (49) can be solved conveniently by existing solvers, such as CVX (24). It is noted that the sufficient and necessary condition for the equivalence of problems P3 – 21’ – 1 and P3 – 21’ with given \( w_2 \), is that the optimal \( W_1^{k+1} \) obtained at k + 1-th iteration of P3 – 21’ – 1 is rank-one, i.e., \( W_1^{k+1} = w_1^{k+1} (w_1^{k+1})^H \), which can be guaranteed by the following lemma.

Lemma 2: \( [23] \) Theorem 2.2] Let \( A_i \in \mathbb{C}^{n \times n}, i \in I = \{1, 2, 3\} \) be a Hermitian matrix and \( X \in \mathbb{H}_+^n \) be a nonzero
Hermitian positive semidefinite matrix. If rank(X) ≥ 2, we can find a rank-one matrix xx† in polynomial-time such that tr(A, xx†) = tr(A, X), i ∈ I.

Proof: The proof of Lemma 2 can be referred to [25].

After achieving the rank-one solution w^{s,k+1} from W^{s,k+1} by using Lemma 2, if P3−21′ is feasible, we plug it into P3−21′. In a similar way to P3−21′−1, for given w^{s,k+1}, we solve the following problem

\[
P3−21′−2: \min_{W^{22}} \text{tr}(W^{22}h_{21}h_{21}^†) \quad (50a)
\]

\[
\text{s.t.:tr}(W^{22}h_{21}h_{21}^†) \geq \begin{cases} t^2-1, & |h_{21}w_{21}^{s,k+1}|^2 + \sigma^2 \\ w_{21}^{s,k+1} & \text{tr}(W^{22}) \leq P_2 \end{cases} \quad (50b)
\]

Similarly, we can obtain the rank-one solution w^{s,k+1} from W^{s,k+1} at k + 1-th iteration based on Lemma 2. If problem P3−21′−2 is feasible. Consequently, we can obtain the optimal solution (w^{s,k+1}, w^{s,k+1}) of P3−21′. Then, by using (w^{s,k+1}, w^{s,k+1}), we can check the feasibility of the constraints (47b) and (47c) of P3−21, which are not considered to obtain (w^{s,k+1}, w^{s,k+1}). If the constraints are feasible, the obtained solution (w^{s,k+1}, w^{s,k+1}) can be a candidate of the optimal beamformers for SIMO case, which are chosen among the solutions of four subproblems. Otherwise, the optimal solution is obtained by solving the other subproblems. The proposed iterative algorithm to solve P3−21 is summarized in the TABLE I.

D. MIMO case (N_1 ≥ 2 and N_2 ≥ 2)

In this subsection, we consider that both Tu_1 and Tu_2 are equipped with N_1 ≥ 2 and N_2 ≥ 2 antennas and hence, the MIMO channel is formed between Tu_1 and Tu_2. For the MIMO case, we jointly design the transmit beamformer at Tu_1, w_1, and beamformers at Tu_2, w_{21} and w_{22}, to maximize the expected achievable rate at Ru_1. In this case, the relaying transmission of Tu_2 cannot improve the achievable rate if the channel quality between Tu_1 and Tu_2 is worse than the direct channel from Tu_1 to Ru_1 such as

\[
\max_{w_1} \log_2 \left(1 + \frac{|w_1^\dagger H_0 w_1|}{\sigma^2} \right) = \log_2 \left(1 + \rho_1 \lambda_{\text{max}}(H_0^\dagger H_0) \right) \leq \log_2 \left(1 + \rho_{\tilde{g}_1} \right), \quad (51)
\]

where \(\lambda_{\text{max}}(X)\) is the largest eigenvalue of X. Therefore, for the MIMO case, the expected achievable rate at Ru_1 is given by

\[
R_{\text{Du_1}}(w_1,w_{21},w_{22}) = \tilde{R}_{\text{Du_1}}(w_1,w_{21},w_{22}) \text{ if } \lambda_{\text{max}}(H_0^\dagger H_0) > \tilde{g}_1, \quad (52)
\]

where \(\tilde{R}_{\text{Du_1}}(w_1,w_{21},w_{22})\) is given by

\[
\tilde{R}_{\text{Du_1}}(w_1,w_{21},w_{22}) = \alpha \text{min} \left\{ \log_2 \left(1 + \frac{|h_{21} w_{21}|^2}{\sigma^2} \right), \log_2 \left(1 + \frac{||H_0 w_1||^2}{\sigma^2} \right) \right\}, \quad (53)
\]

We define the first term in (53) as

\[
Q_{\text{Du_1}}(w_1,w_{21},w_{22}) = \text{min} \left\{ \log_2 \left(1 + \frac{|h_{21} w_{21}|^2}{\sigma^2} \right), \log_2 \left(1 + \frac{||H_0 w_1||^2}{\sigma^2} \right) \right\}. \quad (54)
\]

For the case where Tu_2 helps the transmission of Tu_1, if the rate achieved at Ru_2 in t_1, \(R_{12}(w_1)\), is greater than \(Q_{\text{Du_1}}(w_1,w_{21},w_{22})\), i.e., \(R_{12}(w_1) = \frac{1}{2} \log_2 \left(1 + \frac{|h_{21} w_{21}|^2}{\sigma^2} \right) \geq Q_{\text{Du_1}}(w_1,w_{21},w_{22})\), Ru_2 can decode \(x_1\) in t_1 and thus, Ru_2 can employ SIC to eliminate the effect of \(x_1\) in its received signal in t_2. Therefore, the achievable rate at Ru_2 is given by

\[
R_{\text{Ru_2}}(w_{21},w_{22}) = \begin{cases} R_{\text{SCC}}(w_{22}), & \text{if } R_{12}(w_1) \geq Q_{\text{Du_1}}(w_1,w_{21},w_{22}) \\ R_{\text{DIC}}(w_{21},w_{22}), & \text{if } R_{12}(w_1) < Q_{\text{Du_1}}(w_1,w_{21},w_{22}) \end{cases}, \quad (55)
\]

where \(\tilde{g}_2 = ||h_2||^2, R_{\text{SCC}}(w_{22})\) and \(R_{\text{DIC}}(w_{21},w_{22})\) are given in (42).

For the MIMO case, in order to maximize the expected achievable rate at Ru_1 while guaranteeing QoS requirement at Ru_2, the beamformer at Tu_1, w_1, and the beamformers at Tu_2, w_{21}, and w_{22}, are jointly optimized by the following problem

\[
P4 : \max_{w_1,w_{21},w_{22}} R_{\text{Du_1}}(w_1,w_{21},w_{22}) \quad (56a)
\]

\[
s.t. \quad R_{\text{Ru_2}}(w_{21},w_{22}) \geq Q, \quad (56b)
\]

\[
w_1^\dagger w_1 \leq P_1, w_{21}^\dagger w_{21} + w_{22}^\dagger w_{22} \leq P_2. \quad (56c)
\]
where \( Q \in [0, Q_{\text{max}}] \). Here, when \( Q = Q_{\text{max}} \), \( T_{u_2} \) does not have the residual power to help \( T_{u_1} \) and hence, we have \( w_{21}^\dagger w_{21} = 0 \) and \( w_{22}^\dagger w_{22} = P_2 \). The power allocation at \( T_{u_2} \) is embedded in the beamformer design.

Similar with SIMO case, \( P_4 \) in (56) can be divided into four subproblems with respect to whether \( R_{u_2} \) applies SIC, and the forms of \( R_{u_2}(w_1, w_{21}, w_{22}) \). Due to the space limitation and the similarities of the subproblems in MIMO case, in the following, we consider the subproblem \( P_4 - 21 \), where \( R_{u_2} \) does not apply SIC in \( t_2 \) and \( Q_{T_{u_1}}(w_1, w_{21}, w_{22}) \) is determined by

\[
Q_{T_{u_1}}(w_1, w_{21}, w_{22}) = \frac{1}{2} \log_2 \left( 1 + \frac{|h_{11}^\dagger w_1|^2}{\sigma^2} + \frac{|h_{21}^\dagger w_{21}|^2}{|h_{22}^\dagger w_{22}|^2 + \sigma^2} \right). \tag{57}
\]

Therefore, for this case, the subproblem \( P_3 - 21 \) is represented by

\[
P_{4-21}: \max_{w_1, w_{21}, w_{22}} \log_2 \left( 1 + \frac{|h_{11}^\dagger w_1|^2}{\sigma^2} + \frac{|h_{21}^\dagger w_{21}|^2}{|h_{22}^\dagger w_{22}|^2 + \sigma^2} \right) + (1 - \alpha) \log_2 \left( 1 + \frac{|h_{21}^\dagger w_{21}|^2}{|h_{22}^\dagger w_{22}|^2 + \sigma^2} \right) \tag{58a}
\]

\[
s.t. \log_2 \left( 1 + \frac{|h_{11}^\dagger w_1|^2}{\sigma^2} \right) < \log_2 \left( 1 + \frac{|h_{11}^\dagger w_1|^2}{\sigma^2} + \frac{|h_{21}^\dagger w_{21}|^2}{|h_{22}^\dagger w_{22}|^2 + \sigma^2} \right), \tag{58b}
\]

\[
\log_2 \left( 1 + \frac{|h_{11}^\dagger w_1|^2}{\sigma^2} \right) < \log_2 \left( 1 + \frac{|h_{11}^\dagger w_1|^2}{\sigma^2} + \frac{|h_{21}^\dagger w_{21}|^2}{|h_{22}^\dagger w_{22}|^2 + \sigma^2} \right), \tag{58c}
\]

\[
\frac{1}{2} \log_2 \left( 1 + \frac{|h_{21}^\dagger w_{21}|^2}{|h_{22}^\dagger w_{22}|^2 + \sigma^2} \right) \geq Q, \tag{58d}
\]

\[
w_1^\dagger w_1 \leq P_1, \quad w_{21}^\dagger w_{21} + w_{22}^\dagger w_{22} \leq P_2. \tag{58e}
\]

where the constraint (58d) is the condition that \( R_{u_2} \) cannot apply SIC in \( t_2 \) and the constraint (58e) is the condition that (54) becomes (57). Since \( P_4 - 21 \) is non-convex and \( w_1, w_{21} \) and \( w_{22} \) are coupled in the constraints, it is hard to directly obtain the solution of \( P_4 - 21 \).

In order to obtain the beamformers for MIMO case, we combine the optimal beamformer structure of MISO case, which is given in Lemma 1 to obtain \( w_1 \), and the BCU based iterative algorithm for SIMO case, which is given in TABLE I, to obtain \( w_{21} \) and \( w_{22} \). At one iteration, for given \( w_{21} \) and \( w_{22} \), we obtain \( w_1 \) based on the optimal structure in (29). Then, at the next iteration, \( w_{21} \) and \( w_{22} \) are optimized for the fixed \( w_1 \) by using BCU based algorithm to maximize the expected achievable rate.

For given \( (w_1, w_{21}, w_{22}) \), the terms related to \( (w_{21}, w_{22}) \) can be regarded as the constant and then, we can see that the problem \( P_4 - 21 \) has the same structure of the problem for MISO case, which optimizes \( w_1 \) only. Hence, for given \( (w_{21}, w_{22}) \), we design \( w_1 \) based on the optimal structure of the beamformer for MISO, given in (29). The optimal beamformer for MISO is constructed by two bases \( w_0 = \Pi_{h_0} h_1 / \|\Pi_{h_0} h_1\| \) and \( w_0^\dagger = \Pi_{h_0} h_1^\dagger / \|\Pi_{h_0} h_1^\dagger\| \). Here, since \( h_0 = \theta h_0 \), for some scalar \( \theta \) and \( h_1 = \Pi_{h_0} h_1 + \Pi_{h_0} h_1 \), the optimal beamformer for MISO in (29) can be represented by two bases \( w_0^\text{opt} = \frac{h_0}{\|h_0\|} \) and \( w_1^\text{opt} = \frac{h_1}{\|h_1\|} \) such as

\[
w_1^\text{opt} = \lambda_1 w_0^\text{opt} + \lambda_2 w_1^\text{opt}, \tag{59}
\]

where \( \lambda_1 \) and \( \lambda_2 \) are determined to satisfy \( \|w_1^\text{opt}\|^2 = 1 \). From (59), we note that the optimal structure of the beamformer for MISO is the linear combination of MRT beamformers of channels \( h_0 \) and \( h_1 \). Thus, for MIMO case, we design \( w_1 \) based on the beamformer structure in (59) as

\[
w_1(\lambda) = \sqrt{P_1} \frac{\lambda w_{0}^{\text{opt}} + (1 - \lambda) w_{1}^{\text{opt}}}{\|\lambda w_{0}^{\text{opt}} + (1 - \lambda) w_{1}^{\text{opt}}\|}, \tag{60}
\]

where \( w_0^{\text{opt}} \) is an eigenvector corresponding the largest eigenvalue of \( H_0^\dagger H_0 \) and \( \lambda \) is a real value in \( 0 \leq \lambda \leq 1 \). For the MIMO channel, the eigenvector corresponding the largest eigenvalue of \( H_0^\dagger H_0 \) is the beamformer that maximizes the achievable rate of \( H_0 \), similar to MRT beamformer for MISO channel. Based on (60), for given \( (w_{21}, w_{22}) \), we optimize the coefficient \( \lambda \) to maximize the expected achievable rate at \( R_{u_1} \) by one-dimensional line search.

Similarly, for given \( w_1 \), the terms related to \( w_1 \) can be regarded as the constant and thus, the problem \( P_4 - 21 \) becomes the problem that has the same structure of that for SIMO case, which optimizes \( w_{21} \) and \( w_{22} \) only. Therefore, for given \( w_1 \), we can obtain \( w_{21} \) and \( w_{22} \) by using the BCU based iterative algorithm for SIMO, proposed in Section III-C. Consequently, for MIMO case, we obtain the beamformers \( w_1, w_{21} \) and \( w_{22} \) iteratively to maximize the expected achievable rate at \( R_{u_1} \) while guaranteeing the QoS of \( R_{u_2} \). The details of the proposed algorithm for MIMO case are summarized in TABLE II.

### IV. Simulation Results

In this section, we evaluate the performance of the trust degree based user cooperation for three cases: SISO case where all users have a single antenna (\( N_1 = N_2 = 1 \)), MISO case where \( T_{u_1} \) has the multiple antennas (\( N_1 = 2, N_2 = 1 \)), SIMO case where \( T_{u_2} \) has the multiple antennas (\( N_1 = 1, N_2 = 2 \)) and MIMO case where both \( T_{u_1} \) and \( T_{u_2} \) have the multiple antennas (\( N_1 = N_2 = 2 \)). Unless otherwise specified, we use the average gains of channel elements as \( \{\sigma_{h_1}, \sigma_{h_2}, \sigma_{h_{12}}, \sigma_{h_{21}}\} \) = \{−35, −45, −30, −25, −25\} dB and the expected achievable rates are averaged over 10^4 channel realizations.
for cooperation at Tu1, where $\rho_1 = \rho_2 = 40$ dB and $N_1 = N_2 = 1$.

A. SISO case ($N_1 = N_2 = 1$)

In Fig. 2 for SISO case, we plot the expected achievable rates of Ru1 according to the QoS requirement at Ru2, $Q$, when the transmit SNR at Tu1 and Tu2 are given by $\rho_1 = \rho_2 = 40$ dB. To compare with the proposed user cooperation scheme, which applies SIC, we also plot the expected achievable rates of the user cooperation without SIC and no cooperation ($\alpha = 0$). For the proposed user cooperation, the optimal power allocation at Tu2 for cooperation is obtained as $\beta^\star$ in Theorem 1. In Fig. 2 we show that the expected achievable rate can be significantly increased by the user cooperation when the trust degree between users is high such as $\alpha = 1$. For the case with $\alpha = 0.5$, Tu2 always helps the transmission of Tu1 while when $\alpha = 0.5$, Tu2 helps the transmission of Tu1 with probability 0.5 even if Tu2 has sufficient power budget after achieving QoS. In addition, when the QoS requirement at Ru2 is small, Tu1 can achieve very high expected achievable rate because Tu2 has a large amount of residual power after achieving its QoS and helps the transmission of Tu1 by using large power. By applying SIC at Ru2, the expected achievable rate at Ru1 can be further improved since Tu2 can achieve QoS requirement with small power, and hence Tu2 can allocate more power for cooperation than that without SIC.

In Fig. 3, we show that the expected achievable rate can be increased by the user cooperation when the QoS requirement at Tu2, $Q$, is small or the amount of transmit power budget at Tu2 is large.

B. MISO case ($N_1 = 2$, $N_2 = 1$)

In the MISO case, we jointly design the beamformer at Tu1, $\mathbf{w}_1$, and the power allocation for cooperation at Tu2, $\beta$, according to trust degree, $\alpha$. For the proposed user cooperation, the beamformer at Tu1 is obtained by Theorem 2 and the corresponding power allocation is obtained by one dimensional search. For the comparison, we show the performance of the case that Tu1 simply uses a MRT beamformer, $\mathbf{w}_1 = \mathbf{w}_1^{\text{MRT}} = \frac{\mathbf{h}_1}{\|\mathbf{h}_1\|}$, and Ru2 does not apply SIC. The performance of the no cooperation between users ($\alpha = 0$) with MRT beamformer is also given as a baseline.

In Fig. 5, we show that the optimal power allocation for cooperative relaying at Tu2, $\beta^\omega$, is plotted as a function of the channel gain from Tu2 to Tu1, $g_{21}$, when $\alpha = 0.5$ and $\rho_1 = 40$ dB. In this figure, we use the channel gains as $\{\sigma_{h_{21}}^2, \sigma_{h_{21}}^2, \sigma_{h_{21}}^2, \sigma_{h_{21}}^2\} = \{-32, -40, -30, -32\}$ dB. When $g_{21}$ is weak, from Fig. 3, Tu2 can apply SIC and Tu2 can allocate more power for cooperation. Hence, the optimal power allocation is obtained by $\beta^\star = \beta_{Q_1}$. In contrast, when $g_{21}$ is strong, Tu2 cannot allocate large power for cooperation and the optimal power allocation is obtained by $\beta^\star = \beta_{Q_2}$. Otherwise, Tu2 reduces $\beta$ to help applying SIC at Ru1, and the optimal power allocation is determined by $\beta^\star = \beta_1$, which is a decreasing function of $g_{21}$. In addition, we can see that Tu2 can allocate more power for cooperation when the QoS requirement at Ru2, $Q$, is small or the amount of transmit power budget at Tu2 is large.
In Fig. 4 for the proposed and reference schemes, we plot the expected achievable rates of Ru₁ according to trust degree, α when \( \rho₁ = \rho₂ = 50\text{dB} \). With the growth of α, for the user cooperation schemes, the expected achievable rates of Ru₁ are increased by the cooperation of Tu₂. For the proposed user cooperation, since the beamformer is efficiently designed by considering both trust degree and physical channel qualities, the performance improvement becomes significantly large according to α. However, when Tu₁ uses MRT beamformer, the expected achievable rate improvement is marginal because beamformer is designed independently from α and hence, the benefit from the cooperation cannot be fully exploited even the trust degree is high.

Fig. 5 the optimal power allocation for cooperation at Tu₂, \( \beta^* \), is plotted according to α. When the Tu₁ transmits its data with MRT beamformer, which is independently designed from α, the corresponding power allocation to maximize the expected achievable rate is also determined independently as a constant. For the proposed beamforming, in Fig. 5 we can see that \( \beta^* \) increases with the growth of trust degree α. In the proposed beamforming, when α is high, the direction of beamformer is steered from h₁ to h₀ to fully utilize the cooperation of Tu₂. Due to constraint of DF relaying, the expected achievable rate in (24) is maximized by balancing the minimum rates achieved at Ru₁ (first term) and Tu₂ (second term). Hence, for high α, (24) is maximized by increasing the second term from beamforming design and compensating the first term by assigning large power for cooperation at Tu₂, i.e., \( \beta^* \) is high.

C. SIMO case (\( N₁ = 1, N₂ = 2 \)) and MIMO case (\( N₁ = 2, N₂ = 2 \))

For the SIMO case, we evaluate the performance of the proposed user cooperation based on trust degree in terms of the expected achievable rate. The expected achievable rate at the proposed scheme is achieved by beamformers obtained from the proposed algorithm, which is given in TABLE I. In Fig. 6 for \( Q = 0.5 \) and 1bps/Hz, we plot the expected achievable rates of the proposed and reference schemes versus the trust degree, α when \( \rho₁ = \rho₂ = 50\text{dB} \). Similar to previous subsections, we can see that for high α, the expected achievable rate at Ru₁ is significantly increased by the cooperation with Tu₂ and the performance is further improved by efficiently designing beamformers based on the proposed algorithm. For the MIMO case, we can see the similar phenomena for \( Q = 1 \) and 2bps/Hz in Fig. 7.

In Fig. 7 for \( Q = 0.5 \) and 1bps/Hz, the expected achievable rate at Ru₁ is plotted according to the relaying channel quality from Tu₂ to Ru₁, \( g₂₁ = ||h₂₁||^2 \). In this figure, the trust degree and transmit SNRs are \( \alpha = 0.5 \) and \( \rho₁ = \rho₂ = 50\text{dB} \). From Fig. 7 we first observe that when the gain of relaying channel increases from \(-50\text{dB} \) to \(-25\text{dB} \), the expected achievable rate at Ru₁ can be increased by cooperative transmission from Tu₂ via h₂₁. However, the expected achievable rate does not increase and is saturated when \( g₂₁ \) increases in the regime of \( g₂₁ > -25\text{dB} \). Since when the quality of the relaying channel, h₂₁, is much better than that of channel between Tu₁ and Tu₂, h₀, due to the DF relaying constraint, the rate achieved at Tu₂, which is the second term of (41), is always lower than the rate achieved at Ru₁, which is the first term of (41) for all feasible beamformers. Hence, the expected achievable rate cannot increase and is saturated even if the quality of the relaying channel is sufficiently good. For the MIMO case, we can see the similar phenomena for \( Q = 1 \) and 2bps/Hz in Fig. 8.
V. Conclusion

In this paper, we propose the user cooperation techniques in the multiple antenna system with two communication pairs, i.e., Tu1–Ru1 and Tu2–Ru2, where Tu2 can help the transmission of Tu1 according to the trust degree. For different antenna configurations at Tu1 and Tu2, we design the user cooperation strategies by taking into account the trust degree information as well as channel information. For the SISO case, as a special case, we first propose an optimal power allocation strategy at Tu2, which maximizes the expected achievable rate at Ru1 while guaranteeing QoS requirement at Ru2, according to the channel qualities and QoS requirement. For the MISO case, we provide an optimal structure of beamformer as a linear combination of the weighted channel vectors. Then, based on the optimal structure, we obtain the beamformer that maximizes an approximated expected achievable rate as a function of the trust degree and corresponding power allocation at Tu2. For the SIMO case, to jointly optimize the beamformers of Tu2, we utilize semidefinite relaxation (SDR) technique and block coordinate update (BCU) method to solve the considered problem, and guarantee the rank-one solutions at each step. Furthermore, for the MIMO case, the similarities among problem structures related to MISO, SIMO and MIMO cases inspire us to combine the design of beamformer at Tu1 from MISO and the alternative algorithm from SIMO together to jointly optimize the beamformers at Tu1 and Tu2 to maximize the expected achievable rate at Ru1. Finally, we show that the trust degree between users can be used to significantly improve the expected achievable rate in the user cooperation networks.

Appendix A

Proof of Theorem 1

From (11), we note that the expected achievable rate at Ru1, \( R_{Ru1}(\beta) \), is an increasing function with \( \beta \). Hence, the optimal power allocation is determined by the maximum \( \beta \) that satisfies the QoS requirement at Ru2 such as \( R_{Ru2}(\beta) \geq Q \). If SIC can be applied at Ru2, \( R_{Ru2}(\beta) \) is given by \( R_{Ru2}^{\text{SIC}}(\beta) \) and otherwise, \( R_{Ru2}(\beta) \) is given by \( R_{Ru2}^{\text{NSIC}}(\beta) \). In order to find maximum \( \beta \) that satisfies \( R_{Ru2}(\beta) \geq Q \), we first find the conditions that SIC can be applied at Ru2 and for these conditions, we find the optimal \( \beta \).

From (13), \( Q_{Tu1}(\beta) \) can be represented by two cases of \( \beta_0 \leq \beta \leq 1 \) and \( 0 \leq \beta \leq \beta_0 \). First, for the case of \( \beta_0 \leq \beta \leq 1 \), \( Q_{Tu1}(\beta) \) is given by

\[
Q_{Tu1} = \frac{1}{2} \log_2 (1 + \rho_1 g_0) , \tag{A.1}
\]

which is a constant and independent from \( \beta \). For this case, if \( g_{12} \geq g_0 \), we have \( R_{12} \geq Q_{Tu1} \), and thus, Ru2 can apply SIC. To satisfy QoS requirement at Ru2, we obtain the condition of \( \beta \) as

\[
R_{Ru2}^{\text{SIC}}(\beta) \geq Q \Rightarrow \beta \leq \beta_1, \tag{A.2}
\]

where \( \beta_1 \) is given in (18). In this case, the optimal \( \beta \) is feasible as \( \beta^* = \beta_1 \) if \( \beta_0 \leq \beta_1 \), and thus, the condition of \( Q \) that makes \( \beta^* \) feasible can be obtained by

\[
\beta_0 \leq \beta_1 \Rightarrow Q \geq r_1, \tag{A.3}
\]

where \( r_1 \) is given in (21). Therefore, if \( g_{12} \geq g_0 \) and \( Q \geq r_1 \), the optimal \( \beta \) is obtained by \( \beta^* = \beta_1 \), and QoS of Ru2 is achieved by applying SIC at Ru2.

For the case of \( 0 \leq \beta \leq \beta_0 \), \( Q_{Tu1}(\beta) \) is given by

\[
Q_{Tu1}(\beta) = \frac{1}{2} \log_2 \left( 1 + \rho_1 g_1 + \frac{\beta_2 g_2 g_1}{1 - \beta} \right). \tag{A.4}
\]

For this case, by using (A.4), the condition that can apply SIC at Ru2 is obtained by

\[
R_{12} = \frac{1}{2} \log_2 (1 + \rho_1 g_{12}) \geq Q_{Tu1}(\beta) = g_{12} \geq g_1, \beta \leq \tilde{\beta}_1, \tag{A.5}
\]

where \( \tilde{\beta}_1 \) is given in (20). Thus, if \( \beta_1 \leq \tilde{\beta}_1 \), the optimal \( \beta \) is obtained by \( \beta^* = \beta_1 \), because \( \beta_1 \) is the maximum \( \beta \) that satisfies the QoS requirement at Ru2 as given in (A.2). We obtain the condition of \( Q \) that satisfies \( \beta_1 \leq \beta_0 \) and \( \beta_0 \leq \beta_1 \) as

\[
\begin{align*}
\beta_1 \leq \beta_0 & \Rightarrow Q \geq r_1, \\
\beta_1 \leq \beta_1 & \Rightarrow Q \geq r_2,
\end{align*} \tag{A.6}
\]

where \( r_1 \) is given in (21). Therefore, if \( g_{12} \geq g_1 \) and \( Q \geq \max(r_1, r_2) \), the optimal \( \beta \) is obtained by \( \beta^* = \beta_1 \), and QoS requirement at Ru2 is achieved by applying SIC at Ru2.

On the other hand, if \( \beta_1 \leq \beta_1 \), equivalent \( Q \leq r_2 \), we cannot guarantee the QoS requirement at Ru2 by \( \beta = \beta_1 \), because Ru2 cannot apply SIC when \( \beta = \beta_1 \). Hence, for \( Q \leq r_2 \), the QoS requirement at Ru2 can be guaranteed by \( \beta = \beta_1 \) with applying SIC or \( \beta = \beta_2 \) without applying SIC. Here, \( \beta_2 \) is obtained to satisfy QoS requirement at Ru2 without using SIC as \( R_{Ru2}^{\text{NSIC}}(\beta_2) \geq Q \).

For \( Q \leq r_2 \), since if \( \beta_0 > \beta_1 > \beta_2 \), the optimal \( \beta \) to guarantee QoS requirement at Ru2 is given by \( \beta^* = \beta_1 \), the condition that makes \( \beta^* = \beta_1 \) is obtained as

\[
\begin{align*}
\tilde{\beta}_1 < \beta_0 & \Rightarrow g_0 > g_1, \\
\tilde{\beta}_1 > \beta_2 & \Rightarrow Q > r_3,
\end{align*} \tag{A.7}
\]

where \( r_3 \) is given in (21). Therefore, if \( g_0 > g_{12} \geq g_1 \) and \( r_2 \geq Q > r_3 \), the optimal \( \beta \) is obtained by \( \beta^* = \beta_1 \) and QoS requirement at Ru2 is achieved by applying SIC at Ru2.

Otherwise, the optimal \( \beta \) is obtained by \( \beta^* = \beta_2 \), and QoS requirement at Ru2 is achieved without applying SIC at Ru2.

Fig. 9. The expected achievable rate at Ru1 versus channel gain \( \| h_{21} \|^2 \) with different given QoS, where \( \rho_1 = \rho_2 = 50 \text{dB}, N_1 = N_2 = 2 \) and \( \alpha = 0.5 \).
APPENDIX B
PROOF OF THEOREM 3
For given β, by substituting \(w_1^{\text{opt}}\) of Lemma 1 into (28), \(\hat{R}_{\text{D2D}}(w_1)\) can be rewritten by
\[
\hat{R}_{\text{D2D}}(\eta) = \frac{\alpha}{2} \log_2 \left\{ \rho_1 \min \left\{ g(\eta) + m_1(\beta), f(\eta) \right\} \right\} + \frac{1 - \alpha}{2} \log_2 (\rho_1 g(\eta)),
\] (B.1)
where \(g(\eta)\) and \(f(\eta)\) are given by
\[
g(\eta) \triangleq (\sqrt{\eta} v_1^2 + \sqrt{1 - \eta}) v_2^2 = |h_1^1 w_1|^2,
\] (B.2)
\[
f(\eta) \triangleq \tilde{\eta} g_0 = |h_1^1 w_1|^2.
\] (B.3)
First, if \(\tilde{g}_0\) is large as \(g_0 > (v_1 + v_2)^2\) and power allocation for cooperation, \(\beta\), is small as \(\beta < \beta_{0}\), we have \(f(\eta) > g(\eta) + m_1(\beta)\) for any \(\eta \in \left[\frac{v_1}{v_1 + v_2}\right] \leq \eta \leq 1\) and thus, \(\hat{R}_{\text{D2D}}(\eta)\) in (B.1) can be rewritten as
\[
\hat{R}_{\text{D2D}}(\eta) = \frac{\alpha}{2} \log_2 \left\{ \rho_1 \left( g(\eta) + m_1(\beta) \right) \right\} + \frac{1 - \alpha}{2} \log_2 (\rho_1 g(\eta)).
\] (B.4)
Since (B.4) is a decreasing function with \(\eta\) with \(\eta \in \left[\frac{v_1}{v_1 + v_2}\right] \leq \eta \leq 1\), we can obtain \(\eta^*\) that maximize \(\hat{R}_{\text{D2D}}(\eta)\) as \(\eta^* = \frac{v_1}{v_1 + v_2}\),
Contrarily, if \(\tilde{g}_0\) is small as \(g_0 < v_1\) or \(\beta\) is large as \(\beta > 1\) for \(\eta \geq v_1\), we have \(f(\eta) < g(\eta) + m_1(\beta)\) for any \(\eta \in \left[\frac{v_1}{v_1 + v_2}\right] \leq \eta \leq 1\). In this case, (B.1) can be rewritten as
\[
\hat{R}_{\text{D2D}}(\eta) = \frac{\alpha}{2} \log_2 \left\{ \rho_1 f(\eta) \right\} + \frac{1 - \alpha}{2} \log_2 (\rho_1 g(\eta)).
\] (B.5)
Since (B.5) is a concave function with respect to \(\eta\), we obtain \(\eta^*\) to maximize (B.5) by solving \(\frac{\partial \hat{R}_{\text{D2D}}(\eta)}{\partial \eta} = 0\) as \(\eta^* = \eta_{\text{MISO}}\), which is given in (30).
Otherwise, according to \(\beta\), \(\hat{R}_{\text{D2D}}(\eta)\) in (B.1) can be represented by either (B.4) or (B.5). We first derive \(\eta_{\beta}(\beta)\), which is given in (32), to satisfy \(f(\eta_{\beta}(\beta)) = g(\eta_{\beta}(\beta)) + m_1(\beta)\). Then, for \(\eta_{\beta}(\beta) \leq \eta \leq 1\), \(\hat{R}_{\text{D2D}}(\eta)\) is represented by (B.4), which is a decreasing function of \(\eta\), and thus, we can obtain \(\eta^* = \eta_{\beta}(\beta)\). For \(\eta_{\beta}(\beta) \leq \eta \leq \eta_{\beta}\), (B.1) is rewritten as
\[
\hat{R}_{\text{D2D}}(\eta) = \frac{\alpha}{2} \log_2 \left\{ \rho_1 f(\eta) \right\} + \frac{1 - \alpha}{2} \log_2 (\rho_1 g(\eta)).
\] (B.6)
Since (B.6) is an increasing function with \(\eta\) for \(\eta_{\beta}(\beta) \leq \eta \leq \eta_{\beta}\), we can obtain \(\eta^* = \eta_{\beta}(\beta)\). Otherwise, if \(\eta_{\beta} \leq \eta_{\beta}(\beta)\), we obtain \(\eta^* = \eta_{\beta}(\beta)\) because (B.5) is an increasing function of \(\eta\) for \(\eta_{\beta}(\beta) \leq \eta \leq \eta_{\beta}(\beta)\). Consequently, if \(\eta_{\beta} \leq \eta_{\beta}(\beta)\), we obtain the optimal \(\eta^* = \eta_{\beta}\), and otherwise, the optimal \(\eta^* = \eta_{\beta}(\beta)\). Therefore, the optimal \(\eta^* = \min\{\eta_{\beta}, \eta_{\beta}(\beta)\}\).
Mingxiong Zhao received the B.S. degree in Electrical Engineering and the Ph.D. degree in Information and Communication Engineering from South China University of Technology (SCUT), Guangzhou, China, in 2011 and 2016, respectively. He was a visiting Ph.D. student at University of Minnesota (UMN), Twin Cities, MN, USA, from 2012 to 2013 and Singapore University of Technology and Design (SUTD), Singapore, from 2015 to 2016, respectively. Since 2016, he has been an Assistant Professor at the School of Software, Yunnan University, Kunming, China. His current research interests are physical layer security, cooperative relay communication, and social aware communication systems.

Jemin Lee (S’06-M’11) received the B.S. (with high honors), M.S., and Ph.D. degrees in Electrical and Electronic Engineering from Yonsei University, Seoul, Korea, in 2004, 2007, and 2010, respectively. She was a Postdoctoral Fellow at the Massachusetts Institute of Technology (MIT), Cambridge, MA from 2010 to 2013, and a Temasek Research Fellow at iTrust, Centre for Research in Cyber Security, Singapore University of Technology and Design (SUTD), Singapore from 2014 to 2016. Currently, she is an Assistant Professor at the Department of Information and Communication Engineering, Daegu Gyeongbuk Institute of Science and Technology (DGIST), Daegu, Korea. Her current research interests include physical layer security, wireless security, heterogeneous networks, and machine-type communication. Dr. Lee is currently an Editor for the IEEE Transactions on Wireless Communications and the IEEE Communications Letters, and served as a Guest Editor for the IEEE Wireless Communications, special issue on LTE in Unlicensed Spectrum, 2016, and the ELSEVIER Physical Communication, special issues on Physical Layer Security in 2016 and Heterogeneous and Small Cell Networks in 2014. She received the IEEE ComSoc Asia-Pacific Outstanding Young Researcher Award in 2014, the Temasek Research Fellowship in 2013, the Chun-Gang Outstanding Research Award in 2011, and the IEEE WCSP Best Paper Award in 2014.

Tony Q.S. Quek (S’98-M’08-SM’12) received the B.E. and M.E. degrees in Electrical and Electronics Engineering from Tokyo Institute of Technology, respectively. At MIT, he earned the Ph.D. in Electrical Engineering and Computer Science. Currently, he is a tenured Associate Professor with the Singapore University of Technology and Design (SUTD). He also serves as the Associate Head of ISTD Pillar and the Deputy Director of the SUTD-ZJU IDEA. His main research interests are the application of mathematical, optimization, and statistical theories to communication, networking, signal processing, and resource allocation problems. Specific current research topics include heterogeneous networks, wireless security, internet-of-things, and big data processing.

Dr. Quek has been actively involved in organizing and chairing sessions, and has served as a member of the Technical Program Committee as well as symposium chairs in a number of international conferences. He is serving as the Workshop Chair for IEEE Globecom in 2017, the Tutorial Chair for the IEEE ICC in 2017, and the Special Session Chair for IEEE SPAWC in 2017. He is currently an elected member of IEEE Signal Processing Society SPCOM Technical Committee. He was an Executive Editorial Committee Member for the IEEE Transactions on Wireless Communications, an Editor for the IEEE Transactions on Communications, and an Editor for the IEEE Wireless Communications Letters. He is a co-author of the book “Small Cell Networks: Deployment, PHY Techniques, and Resource Allocation” published by Cambridge University Press in 2013 and the book “Cloud Radio Access Networks: Principles, Technologies, and Applications” by Cambridge University Press in 2017.

Dr. Quek was honored with the 2008 Philip Yeo Prize for Outstanding Achievement in Research, the IEEE Globecom 2010 Best Paper Award, the 2012 IEEE William R. Bennett Prize, the IEEE SPAWC 2013 Best Student Paper Award, the IEEE WCSP 2014 Best Paper Award, the 2015 SUTD Outstanding Education Awards – Excellence in Research, the 2016 Thomson Reuters Highly Cited Researcher, and the 2016 IEEE Signal Processing Society Young Author Best Paper Award.

Suili Feng (M’05) received the B.S. degree in Electrical Engineering and the M.S. and Ph.D. degrees in Electronic and Communication System from South China University of Technology, Guangzhou, China, in 1982, 1989, and 1998, respectively. He was a Research Assistant with Hong Kong Polytechnic University, Hong Kong, from 1991 to 1992, and a Visiting Scholar with University of South Florida, Tampa, FL, USA, from 1998 to 1999. He has been with South China University of Technology since 1989, where he is currently a Professor with the School of Electronic and Information Engineering. His research interests include wireless networks, computer networks, and communication signal processing.

Jong Yool Ryu (S’11-M’14) received the B.E. degree in Electrical Engineering from Chungnam National University, Daejeon, Korea, in 2008. He received the M.S. and Ph.D. degrees in Electrical Engineering from Korea Advance Institute of Science and Technology (KAIST), Daejeon, Korea, in 2010 and 2014, respectively. From 2014 to 2016, he was a Postdoctoral Fellow at Singapore University of Technology and Design (SUTD). Currently, he is an Assistant Professor at the Department of Information and Communication Engineering, Gyeongsang National University (GNU), Tongyeong, Korea. His current research interests are communication secrecy, cooperative user relay communications, and social aware communication systems.