Fractal Intersections and Products via Algorithmic Dimension

Neil Lutz*
Department of Computer Science, Rutgers University
Piscataway, NJ 08854, USA
njlutz@rutgers.edu

January 18, 2017

Abstract
Algorithmic dimensions quantify the algorithmic information density of individual points and may be defined in terms of Kolmogorov complexity. This work uses these dimensions to bound the classical Hausdorff and packing dimensions of intersections and Cartesian products of fractals in Euclidean spaces. This approach shows that a known intersection formula for Borel sets holds for arbitrary sets, and it significantly simplifies the proof of a known product formula. Both of these formulas are prominent, fundamental results in fractal geometry that are taught in typical undergraduate courses on the subject.

1 Introduction

Classical fractal dimensions, among which Hausdorff dimension \[11\] is the most important, refine notions of measure to quantitatively classify sets of measure 0. In 2000, J. Lutz \[14\] showed that Hausdorff dimension can be simply characterized using betting strategies called gales, and that this characterization can be effectivized in order to quantitatively classify non-random infinite data objects. This effective Hausdorff dimension and other, related algorithmic dimensions have been applied to multiple areas of computer science and have proven especially useful in algorithmic information theory \[10, 25, 5\].

The connection between algorithmic and classical dimensions has more recently been exploited in the other direction, i.e., to apply algorithmic information theoretic methods and intuition to classical fractal geometry (e.g., \[27, 2\]). A point-to-set principle of J. Lutz and N. Lutz \[15\], stated here as Theorem \[7\], characterizes the classical Hausdorff dimension of any set in \(\mathbb{R}^n\) in terms of the algorithmic dimensions of its individual points.

*Research supported in part by National Science Foundation Grant 1445755.
In the same work, J. Lutz and N. Lutz showed that this principle gives rise to a new, pointwise technique for dimensional lower bounds, and, as a proof of concept, used this technique to give an algorithmic information theoretic proof of Davies’s 1971 theorem stating that every Kakeya set in \( \mathbb{R}^2 \) has Hausdorff dimension 2. This bounding technique has since been used by N. Lutz and Stull to make new progress on a problem in classical fractal geometry by deriving an improved lower bound on the Hausdorff dimension of generalized Furstenberg sets, as defined by Molter and Rela.

The same algorithmic dimensional technique is applied here to bound the dimensions of intersections and products of fractals. Most significantly, we extend the following intersection formula, previously shown to hold when \( E \) and \( F \) are Borel sets, to arbitrary sets \( E \) and \( F \).

**Theorem 1.** For all \( E, F \subseteq \mathbb{R}^n \), and for almost every \( z \in \mathbb{R}^n \),

\[
\dim_H(E \cap (F + z)) \leq \max\{0, \dim_H(\mathbb{R}^n) - n\},
\]

where \( F + z = \{x + z : x \in F\} \).

This approach also yields a simplified proof of the following known product formula for general sets.

**Theorem 2 (Marstrand).** For all \( E \subseteq \mathbb{R}^m \) and \( F \subseteq \mathbb{R}^n \),

\[
\dim_H(E) + \dim_H(F) \leq \dim_H(E \times F).
\]

We use symmetric arguments to derive the known corresponding statements about packing dimension, a formulation of fractal dimension that was developed independently by Tricot and Sullivan and is dual to Hausdorff dimension. These results are included here to showcase the versatility of this technique and its ability to capture the exact duality between Hausdorff and packing dimensions.

## 2 Classical Fractal Dimensions

We begin by stating classical, measure-theoretic definitions of the two most well-studied notions of fractal dimension, Hausdorff dimension and packing dimension. These definitions are included here for completeness but are not used directly in the remainder of this work; we will instead apply equivalent characterizations in terms of algorithmic information, as described in Section 3.

**Definition (Hausdorff).** For \( E \subseteq \mathbb{R}^n \), let \( \mathcal{U}_\delta(E) \) be the collection of all countable covers of \( E \) by sets of positive diameter at most \( \delta \), where the **diameter** of any set \( U \subseteq \mathbb{R}^n \) is given by

\[
diam(U) = \sup_{x,y \in U} |x - y|.
\]
For all \( s \geq 0 \), let
\[
H^s_\delta(E) = \inf \left\{ \sum_{i \in \mathbb{N}} \text{diam}(U_i)^s : \{U_i\}_{i \in \mathbb{N}} \in \mathcal{U}_\delta(E) \right\}.
\]
The \textit{s-dimensional Hausdorff (outer) measure} of \( E \) is
\[
H^s(E) = \lim_{\delta \to 0^+} H^s_\delta(E),
\]
and the \textit{Hausdorff dimension} of \( E \) is
\[
\dim_H(E) = \inf \{ s > 0 : H^s(E) = 0 \} = \sup \{ s : H^s(E) = \infty \}.
\]

Three desirable properties have made \( \dim_H \) the most standard notion of fractal dimension since it was introduced by Hausdorff in 1919. First, it is defined on every set in \( \mathbb{R}^n \). Second, it is \textit{monotone}: if \( E \subseteq F \), then \( \dim_H(E) \leq \dim_H(F) \). Third, it is \textit{countably stable}: if \( E = \bigcup_{i \in \mathbb{N}} E_i \), then \( \dim_H(E) = \sup_{i \in \mathbb{N}} \dim_H(E_i) \). These three properties also hold for packing dimension, which was defined much later, independently by Tricot [30] and by Sullivan [29].

**Definition** (Tricot [30], Sullivan [29]). For all \( x \in \mathbb{R}^n \) and \( \rho > 0 \), let \( B_\rho(x) \) denote the open ball of radius \( \rho \) and center \( x \). For all \( E \subseteq \mathbb{R}^n \), let \( \mathcal{V}_\delta(E) \) be the collection of all countable packings of \( E \) by open balls of diameter at most \( \delta \). That is, for every packing \( \{V_i\}_{i \in \mathbb{N}} \in \mathcal{V}_\delta(E) \) and every \( i \in \mathbb{N} \), we have \( V_i = B_{\rho_i}(x_i) \) for some \( x_i \in E \) and \( \rho_i \in [0, \delta/2] \). For all \( s \geq 0 \), define
\[
P^s_\delta(E) = \sup \left\{ \sum_{i \in \mathbb{N}} \text{diam}(V_i)^s : \{V_i\}_{i \in \mathbb{N}} \in \mathcal{V}_\delta(E) \right\},
\]
and let
\[
P^s_0(E) = \lim_{\delta \to 0^+} P^s_\delta(E).
\]
The \textit{s-dimensional packing (outer) measure} of \( E \) is
\[
P^s(E) = \inf \left\{ \sum_{i \in \mathbb{N}} P^s_0(E_i) : E \subseteq \bigcup_{i \in \mathbb{N}} E_i \right\},
\]
and the \textit{packing dimension} of \( E \) is
\[
\dim_P(E) = \inf \{ s > 0 : P^s(E) = 0 \} = \sup \{ s > 0 : P^s(E) = \infty \}.
\]

Notice that defining packing dimension in this way requires an extra step of optimization compared to Hausdorff dimension. More properties and details about classical fractal dimensions may be found in standard references such as [22, 9, 28].
3 Algorithmic Fractal Dimensions

This section defines the effective Hausdorff and packing dimensions in terms of algorithmic information, i.e., Kolmogorov complexity. We also define conditional and mutual dimensions and discuss some properties of these dimensions, including their relationships to classical Hausdorff and packing dimensions.

3.1 Kolmogorov Complexity

Kolmogorov complexity is most often defined in the space \( \{0,1\}^* \) of binary strings, but it is readily extended to other discrete domains. For the purposes of this work, the complexity of rational points is most relevant. Hence, fix some standard binary encoding for \( n \)-tuples of rationals. The Kolmogorov complexity of \( p \) is the length of the shortest binary program that outputs \( p \). Formally, it is

\[
K(p) = \min_{\pi \in \{0,1\}^*} \{|\pi| : U(\pi) = p\},
\]

where \( U \) is a fixed universal prefix-free Turing machine and \(|\pi|\) is the length of \( \pi \). This quantity is also called the algorithmic information content of \( p \).

The conditional Kolmogorov complexity of \( p \) given \( q \in \mathbb{Q}^n \) is the length of the shortest binary program that outputs \( p \) when given \( q \) as an input:

\[
K(p|q) = \min_{\pi \in \{0,1\}^*} \{|\pi| : U(\pi, q) = p\}.
\]

The algorithmic mutual information between \( p \in \mathbb{Q}^m \) and \( q \in \mathbb{Q}^n \) measures, informally, the amount that knowledge of \( q \) helps in the task of compressing \( p \). Formally, it is

\[
I(p : q) = K(p) - K(p|q).
\]

The quantities \( K(p), K(p|q), \) and \( I(p : q) \) are closely related to, and may be considered algorithmic analogues of, the quantities entropy \( H(X) \), conditional entropy \( H(X|Y) \), and mutual information \( I(X;Y) \), from classical (Shannon) information theory. See references \([13, 25, 5]\) for more details on algorithmic information and the connections between algorithmic and classical theories of information.

3.2 Effective Dimensions

Using approximation by rationals, Kolmogorov complexity may be further extended to Euclidean spaces \([16]\). For every \( E \subseteq \mathbb{R}^n \), define

\[
K(E) = \min\{K(p) : p \in E \cap \mathbb{Q}^n\},
\]

where the minimum is understood to be infinite if \( E \cap \mathbb{Q}^n \) is empty. This is the length of the shortest program that outputs some rational point in \( E \). The Kolmogorov complexity of \( x \in \mathbb{R}^n \) at precision \( r \in \mathbb{N} \) is given by

\[
K_r(x) = K(B_2^{n}(x)),
\]
the length of the shortest program that outputs any precision-$r$ rational approximation of $x$. $K_r(x)$ may also be described as the algorithmic information content of $x$ at precision $r$, and similarly, $K_r(x)/r$ is the algorithmic information density or incompressibility of $x$ at precision $r$. This ratio does not necessarily converge as $r \to \infty$, but it does have asymptotes in $[0, n]$. These asymptotes are used to define effective dimensions.

**Definition** ([14, 23, 1, 16]). Let $x \in \mathbb{R}^n$.

1. The effective Hausdorff dimension of $x$ is

   $$\dim(x) = \liminf_{r \to \infty} \frac{K_r(x)}{r}.$$ 

2. The effective packing dimension of $x$ is

   $$\text{Dim}(x) = \limsup_{r \to \infty} \frac{K_r(x)}{r}.$$ 

These dimensions were originally defined by J. Lutz [14] and Athreya, Hitchcock, J. Lutz, and Mayordomo [1], respectively. The original definitions were in Cantor space and used *gales*, which are betting strategies that generalize martingales. Their Kolmogorov complexity characterizations and translation to Euclidean spaces are due to Mayordomo [23] and J. Lutz and Mayordomo [16].

The information theoretic nature of these characterizations has led to the development of algorithmic dimensional quantities corresponding to the other algorithmic information theoretic quantities defined above. As analogues to mutual information and conditional information, Case and J. Lutz defined *mutual dimensions* and J. Lutz and N. Lutz defined *conditional dimensions*.

### 3.3 Mutual Dimensions

Given $E \subseteq \mathbb{R}^m$ and $F \subseteq \mathbb{R}^n$, define

$$I(E : F) = \min \{ I(p : q) : p \in E \cap \mathbb{Q}^m \text{ and } q \in F \cap \mathbb{Q}^n \}.$$ 

Then the mutual information between $x \in \mathbb{R}^m$ at precision $r \in \mathbb{N}$ and $y \in \mathbb{R}^n$ at precision $s \in \mathbb{N}$ is given by

$$I_{r,s}(x : y) = I(B_{2^{-r}}(x) : B_{2^{-s}}(y)).$$ 

**Definition** (Case and J. Lutz [3]). Let $x \in \mathbb{R}^m$ and $y \in \mathbb{R}^n$.

1. The lower mutual dimension between $x$ given $y$ is

   $$\text{mdim}(x : y) = \liminf_{r \to \infty} \frac{I_r(x : y)}{r}.$$
2. The upper mutual dimension between $x$ and $y$ is

$$\text{Mdim}(x : y) = \limsup_{r \to \infty} \frac{I_{r,r}(x : y)}{r}.$$  

**Observation 3.** For all $x \in \mathbb{R}^n$,

1. $\dim(x) = \text{mdim}(x : x)$.
2. $\text{Dim}(x) = \text{Mdim}(x : x)$.

Case and J. Lutz also showed that mutual dimensions are preserved under bi-Lipschitz computable bijections. Combined with Observation 3, this implies that effective Hausdorff and packing dimensions are preserved by such bijections. This is a Euclidean-space version of a fact that was shown in Cantor space by Reimann.

**Lemma 4** (Reimann [20], Case and J. Lutz [3]). If $f : \mathbb{R}^m \to \mathbb{R}^n$ is computable and bi-Lipschitz, then $\dim(x) = \dim(f(x))$ and $\text{Dim}(x) = \text{Dim}(f(x))$ for all $x \in \mathbb{R}^m$.

### 3.4 Conditional Dimensions

Given $E \subseteq \mathbb{R}^m$ and $F \subseteq \mathbb{R}^n$, define

$$K(E|F) = \max \{ \min \{ K(p|q) : p \in E \cap \mathbb{Q}^m \} : q \in F \cap \mathbb{Q}^n \}.$$  

Then the conditional Kolmogorov complexity of $x \in \mathbb{R}^m$ at precision $r \in \mathbb{N}$ given $y \in \mathbb{R}^n$ at precision $s \in \mathbb{N}$ is given by

$$K_{r,s}(x|y) = K(B_{2^{-r}}(x)|B_{2^{-s}}(y)).$$  

**Definition** (J. Lutz and N. Lutz [15]). Let $x \in \mathbb{R}^m$ and $y \in \mathbb{R}^n$.

1. The lower conditional dimension of $x$ given $y$ is

$$\dim(x : y) = \liminf_{r \to \infty} \frac{K_{r,s}(x|y)}{r}.$$  

2. The upper conditional dimension of $x$ given $y$ is

$$\text{Dim}(x : y) = \limsup_{r \to \infty} \frac{K_{r,s}(x|y)}{r}.$$  

That work also showed that the symmetry of algorithmic information holds in Euclidean space, in the form

$$K_r(x, y) = K_{r,r}(x) + K_r(y|x) + o(r).$$  

This fact and elementary properties of limits inferior and superior immediately imply the following chain rule for effective dimensions.
Theorem 5 (J. Lutz and N. Lutz [15]). For all $x \in \mathbb{R}^m$ and $y \in \mathbb{R}^n$,
\[
\dim(x) + \dim(y|x) \leq \dim(x, y) \leq \dim(x) + \dim(y|x) \leq \dim(x) + \dim(y|x).
\]

3.5 Oracles and Relative Dimensions

By making the fixed universal machine $U$ an oracle machine, the algorithmic information quantities above may be defined relative to any oracle $A \subseteq \mathbb{N}$. The definitions of $K^A(\sigma|\tau)$, $K^A(\sigma)$, $K^A(x)$, $K^A(x|y)$, $\dim^A(x)$, $\dim^A(x|y)$ and $\text{Dim}^A(x|y)$ all exactly mirror their unrelativized versions, except that $U$ is permitted to query membership in $A$ as a computational step.

For $y \in \mathbb{R}^n$, we write $\dim^y(x)$ as shorthand for $\dim^{A_y}(x)$, where $A_y \subseteq \mathbb{N}$ encodes the binary expansions of $y$’s coordinates in some standard way, and similarly for $\text{Dim}^y(x)$. Since this kind of oracle access to $y$ is at least as informative as any finite-precision estimate for $y$ (ignoring the small amount of information given by the precision parameter itself), these relative dimensions are bounded above by conditional dimensions.

Lemma 6 (J. Lutz and N. Lutz [15]). For all $x \in \mathbb{R}^m$ and $y \in \mathbb{R}^n$,
\[
1. \dim^y(x) \leq \dim(x|y),
2. \text{Dim}^y(x) \leq \text{Dim}(x|y).
\]

3.6 Point-to-Set Principle

Effective Hausdorff dimension and effective packing dimension were conceived as constructive versions of classical Hausdorff dimension and packing dimension [14, 1]. The following point-to-set principle uses relativization to precisely characterize their relationships to their non-algorithmic precursors.

Theorem 7 (J. Lutz and N. Lutz [15]). For every $E \subseteq \mathbb{R}^n$, the Hausdorff dimension and packing dimension of $E$ are
\[
1. \dim_H(E) = \min_{A \subseteq \mathbb{N}} \sup_{x \in E} \dim^A(x),
2. \dim_P(E) = \min_{A \subseteq \mathbb{N}} \sup_{x \in E} \text{Dim}^A(x).
\]

Notice that, unlike the definitions of $\dim_H(E)$ and $\dim_P(E)$ given in Section 2, the above characterizations are completely symmetrical.

Theorem 7 allows us to prove lower bounds on classical dimensions in a pointwise way. To show a statement of the form $\dim_H(E) \geq \alpha$, it suffices to show, for a given oracle $A$ and every $\varepsilon > 0$, that there exists an $x \in E$ satisfying $\dim^A(x) \geq \alpha - \varepsilon$. This often be argued using Kolmogorov complexity, but the proofs in Sections 4 and 5 do not directly use Kolmogorov complexity; the only tools needed are Lemma 4, Theorem 5, Lemma 6, and Theorem 7.
4 Intersections of Fractals

In this section we prove Theorem 1. We then use a symmetric argument to prove the corresponding statement for packing dimension, which is known \[8\]. For the case where \(E, F \subseteq \mathbb{R}^n\) are Borel sets, Theorem 1 was shown in its present form by Falconer \[7\]. Closely related results, which also place restrictions on \(E\) and \(F\), were proven earlier by Mattila \[20, 21\] and Kahane \[12\].

**Theorem 1.** For all \(E, F \subseteq \mathbb{R}^n\), and for almost every \(z \in \mathbb{R}^n\),

\[
\dim_H(E \cap (F + z)) \leq \max\{0, \dim_H(E \times F) - n\}, \tag{1}
\]

where \(F + z = \{x + z : x \in F\}\).

**Proof.** Let \(E, F \subseteq \mathbb{R}^n\) and \(z \in \mathbb{R}^n\). If \(E \cap (F + z) = \emptyset\), then (1) holds trivially, so assume that the intersection is nonempty. Theorem 7a guarantees that there is some oracle set \(A \subseteq \mathbb{N}\) satisfying

\[
\dim_H(E \times F) = \sup_{(x, y) \in E \times F} \dim_A(x, y). \tag{2}
\]

It also guarantees, given any \(\varepsilon > 0\), that there is an \(x \in E \cap (F + z)\) such that

\[
\dim_A(x, z) \geq \dim_H(E \cap (F + z)) - \varepsilon. \tag{3}
\]

Since \((x, x - z) \in E \times F\), we have

\[
\dim_H(E \times F) \geq \dim_A(x, x - z)
\]

\[= \dim_A(x, z)
\]

\[\geq \dim_A(z) + \dim_A(x|z)
\]

\[\geq \dim_A(z) + \dim_A(x) + \dim_A(z)\]

\[\geq \dim_A(z) + \dim_H(E \cap (F + z)) - \varepsilon.
\]

The above lines follow from Lemma 4, Theorem 5, Lemma 6, and 8, respectively. Letting \(\varepsilon \to 0\), we have

\[
\dim_H(E \cap (F + z)) \leq \dim_H(E \times F) - \dim_A(z).
\]

Thus, (a) holds whenever \(\dim_A(z) = n\). In particular, it holds when \(z\) is Martin-Löf random relative to \(A\), i.e., for Lebesgue almost every \(z \in \mathbb{R}^n\) \[13, 19\].

For the case that \(E\) and \(F\) are Borel sets, Falconer \[9\] notes that the intersection formula is readily extended to rigid motions and similarities. The same argument applies in the general case, so Theorem 1 has the following corollary.

**Corollary 8.** Let \(E, F \subseteq \mathbb{R}^n\). Let \(G\) be the group of rigid motions or the group of similarities on \(\mathbb{R}^n\). Then, for almost all \(\sigma \in G\),

\[
\dim_H(E \cap \sigma(F)) \leq \max\{0, \dim_H(E \times F) - n\}. \tag{4}
\]
Figure 1: Let $E$ and $F$ each be Koch snowflakes, which have Hausdorff dimension $\log_3 4 \approx 1.26$. Left: For almost all rigid motions $\sigma$, the intersection $E \cap \sigma(F)$ has Hausdorff dimension at most $2 \log_3 4 - 2 \approx 0.52$. Right: For a measure zero set of rigid motions, the Hausdorff dimension of the intersection may be as large as $\log_3 4$. Note that Koch curves are Borel sets, so the new generality given by Theorem 1 and Corollary 8 is not required for this example.

Proof (Following Falconer [9]). For all rotations (and all scalings) of $F$, Theorem 1 tells us that holds for almost all translations. Thus, holds for almost all rigid motions and almost all similarities.

A corresponding intersection formula for packing dimension has been shown for arbitrary $E, F \subseteq \mathbb{R}^n$ by Falconer [8]. That proof is not difficult or long, but an algorithmic dimensional proof is presented here as an instance where this technique applies symmetrically to both Hausdorff and packing dimension.

**Theorem 9** (Falconer [8]). For all $E, F \subseteq \mathbb{R}^n$, and for almost every $z \in \mathbb{R}^n$,

$$\dim_P(E \cap (F + z)) \leq \max\{0, \dim_P(E \times F) - n\}.$$  

**Proof.** As in Theorem 1, we may assume that the intersection is nonempty. Apply Theorem 7 to choose an oracle set $B \subseteq \mathbb{N}$ such that

$$\dim_P(E \times F) = \sup_{(x,y) \in E \times F} \Dim_B(x, y)$$  

and, given $\varepsilon > 0$, a point $y \in E \cap (F + z)$ satisfying

$$\Dim_B(z, y) \geq \dim_P(E \cap (F + z)) - \varepsilon.$$  

Then $(y, y - z) \in E \times F$, and we may proceed much as before:

$$\dim_P(E \times F) \geq \Dim_B(y, y - z)$$

$$= \Dim_B(y, z)$$

$$\geq \dim_B(z) + \dim_B(y|z)$$

$$\geq \dim_B(z) + \Dim_B(z, y)$$

$$\geq \dim_B(z) + \dim_P(E \cap (F + z)) - \varepsilon.$$  

These lines follow from [5], Lemma 4, Theorem 5, Lemma 6 and 8. Again, $\dim_B(z) = n$ for almost every $z \in \mathbb{R}^n$, so this completes the proof. 

\[9\]
5 Products of Fractals

In this section we prove four known product inequalities for fractal dimensions. Inequality (7), which was stated in the introduction as Theorem 2, is due to Marstrand [18]. When $E$ and $F$ are Borel sets, it is simple to prove (7) by using Frostman’s Lemma, but the argument for general sets using net measures is considerably more difficult [22, 6]. The other three inequalities are due to Tricot [30]. Reference [22] gives a more detailed account of this history.

Theorem 10 (Marstrand [18], Tricot [30]). For all $E \subseteq \mathbb{R}^m$ and $F \subseteq \mathbb{R}^n$,

\[
\dim_H(E) + \dim_H(F) \leq \dim_H(E \times F) \leq \dim_H(E) + \dim_P(F) \leq \dim_P(E \times F) \leq \dim_P(E) + \dim_P(F).
\]

Notice the superficial resemblance of this theorem to Theorem 5. This similarity is not a coincidence; each inequality in Theorem 10 follows from the corresponding line in Theorem 5. The arguments given here for (7–10) are each similar in length to the proof of (7) for Borel sets. That is, they are quite short.

Proof. Theorem 7a guarantees, for every $\varepsilon > 0$, that there exist an oracle set $A \subseteq \mathbb{N}$ and points $x \in E$ and $y \in F$ such that

\[
\dim_H(E \times F) = \sup_{z \in E \times F} \dim^A(z), \quad \dim^A(x) \geq \dim_H(E) - \varepsilon, \quad \dim^{A,x}(y) \geq \dim_H(F) - \varepsilon.
\]

Then by (11), Theorem 5 relative to $A$, and Lemma 6 relative to $A$, we have

\[
\dim_H(E \times F) \geq \dim^A(x, y) \geq \dim^A(x) + \dim^A(y|x) \geq \dim^A(x) + \dim^{A,x}(y) \geq \dim_H(E) + \dim_H(F) - 2\varepsilon,
\]

by our choice of $x$ and $y$. Since $\varepsilon > 0$ was arbitrary, we conclude that (7) holds.

For (8), let $\varepsilon > 0$ and use both parts of Theorem 7 to find $B, C \subseteq \mathbb{N}$, $u \in E$, and $v \in F$ such that

\[
\dim_H(E) = \sup_{x \in E} \dim^B(x), \quad \dim_P(F) = \sup_{y \in E} \dim^C(y), \quad \dim^B,C(u, v) \geq \dim_H(E \times F) - \varepsilon.
\]
Since $B$ and $C$ minimize their respective expressions, we also have
\[
\dim_H(E) = \sup_{x \in E} \dim^{B,C}(x),
\]
\[
\dim_P(F) = \sup_{y \in E} \Dim^{B,C}(y).
\]
Thus, we can apply Theorem 5 relative to $B, C$, after first noticing that conditioning on another point never increases dimension.

\[
\dim_H(E) + \dim_P(F) \geq \dim^{B,C}(u) + \Dim^{B,C}(v)
\]
\[
\geq \dim^{B,C}(u|v) + \Dim^{B,C}(v)
\]
\[
\geq \dim^{B,C}(u, v)
\]
\[
\geq \dim_H(E \times F) - \varepsilon.
\]

Again, $\varepsilon$ was arbitrary, so (8) holds.

For (9) and (10), we use essentially the same arguments as above. By Theorem 7, there are $A', B' \subseteq \mathbb{N}$, $x', u' \in E$, $y', v' \in F$, and $\varepsilon > 0$ that satisfy

\[
\dim_P(E \times F) = \sup_{z \in E \times F} \Dim^{A'}(z),
\]
\[
\dim_H(E) = \sup_{z \in E} \Dim^{B'}(z),
\]
\[
\dim^{A'}(x') \geq \dim_H(E) - \varepsilon,
\]
\[
\Dim^{A',x'}(y') \geq \dim_P(F) - \varepsilon,
\]
\[
\Dim^{B',C}(u', v') \geq \dim_P(E \times F) - \varepsilon,
\]
where $x$ and $C$ are as above. We once again apply relativized versions of Theorem 5 and Lemma 6:
\[
\dim_P(E) + \dim_P(F) \geq \Dim^{B',C}(u') + \Dim^{B',C}(v')
\]
\[
\geq \Dim^{B',C}(u'|v') + \Dim^{B',C}(v')
\]
\[
\geq \Dim^{B',C}(u', v')
\]
\[
\geq \dim_P(E \times F) - \varepsilon
\]
\[
\geq \Dim^{A'}(x', y') - \varepsilon
\]
\[
\geq \dim^{A'}(x') + \Dim^{A'}(y'|x') - \varepsilon
\]
\[
\geq \dim^{A'}(x') + \Dim^{A',x'}(y') - \varepsilon
\]
\[
\geq \dim_H(E) + \dim_P(F) - 3\varepsilon.
\]

Letting $\varepsilon \to 0$ completes the proof.

\section{Conclusion}

The applications of theoretical computer science to pure mathematics in this paper yielded a significant extension to a basic theorem on Hausdorff dimension,
as well as a much simpler argument for another such theorem. Understanding classical fractal dimensions as pointwise, algorithmic information theoretic quantities enables reasoning about them in a way that is both fine-grained and intuitive, and the proofs in this work are further evidence of the power and versatility of bounding techniques using Theorem 7. In particular, Theorem 1 demonstrates that this approach can be used to strengthen the foundations of fractal geometry. Therefore, in addition to further applications of these techniques, developing more refined results on the relationship between classical geometric measure theory and Kolmogorov complexity is an appealing direction for future investigations.

References

[1] Krishna B. Athreya, John M. Hitchcock, Jack H. Lutz, and Elvira Mayordomo. Effective strong dimension in algorithmic information and computational complexity. *SIAM Journal of Computing*, 37(3):671–705, 2007.

[2] Verónica Becher, Jan Reimann, and Theodore A. Slaman. Irrationality exponent, hausdorff dimension and effectivization. 2016.

[3] Adam Case and Jack H. Lutz. Mutual dimension. *ACM Transactions on Computation Theory*, 7(3):12, 2015.

[4] R. O. Davies. Some remarks on the Kakeya problem. *Proceedings of the Cambridge Philosophical Society*, 69:417–421, 1971.

[5] Rod Downey and Denis Hirschfeldt. *Algorithmic Randomness and Complexity*. Springer-Verlag, 2010.

[6] Kenneth J. Falconer. *The Geometry of Fractal Sets*. Cambridge University Press, 1985.

[7] Kenneth J. Falconer. *Fractal Geometry: Mathematical Foundations and Applications*. Wiley, first edition, 1990.

[8] Kenneth J. Falconer. Sets with large intersection properties. *Journal of the London Mathematical Society*, 49(2):267–280, 1994.

[9] Kenneth J. Falconer. *Fractal Geometry: Mathematical Foundations and Applications*. Wiley, third edition, 2014.

[10] Xiaoyang Gu and Jack H. Lutz. Dimension characterizations of complexity classes. *Computational Complexity*, 17(4):459–474, 2008.

[11] Felix Hausdorff. Dimension und äusseres Mass. *Mathematische Annalen*, 79:157–179, 1919.

[12] Jean-Pierre Kahane. Sur la dimension des intersections. In Jorge Alberto Barroso, editor, *Aspects of mathematics and its applications*, North-Holland Mathematical Library, 34, pages 419–430. Elsevier, 1986.
[13] Ming Li and Paul M.B. Vitányi. An Introduction to Kolmogorov Complexity and Its Applications. Springer, third edition, 2008.

[14] Jack H. Lutz. The dimensions of individual strings and sequences. Information and Computation, 187(1):49–79, 2003.

[15] Jack H. Lutz and Neil Lutz. Algorithmic information, plane Kakeya sets, and conditional dimension. Symposium on Theoretical Aspects of Computer Science, STACS 2017, Hannover, Germany, to appear.

[16] Jack H. Lutz and Elvira Mayordomo. Dimensions of points in self-similar fractals. SIAM Journal of Computing, 38(3):1080–1112, 2008.

[17] Neil Lutz and D. M. Stull. Bounding the dimension of points on a line. Proceedings of the 14th Annual Conference on Theory and Applications of Models of Computation, TAMC 2017, Bern, Switzerland, to appear.

[18] John M. Marstrand. Some fundamental geometrical properties of plane sets of fractional dimensions. Proceedings of the London Mathematical Society, 4(3):257–302, 1954.

[19] Per Martin-Löf. The definition of random sequences. Information and Control, 9(6):602–619, 1966.

[20] Pertti Mattila. Hausdorff dimension and capacities of intersections of sets in n-space. Acta Mathematica, 152:77–105, 1984.

[21] Pertti Mattila. On the Hausdorff dimension and capacities of intersections. Mathematika, 32:213–217, 1985.

[22] Pertti Mattila. Geometry of sets and measures in Euclidean spaces: fractals and rectifiability. Cambridge University Press, 1995.

[23] Elvira Mayordomo. A Kolmogorov complexity characterization of constructive Hausdorff dimension. Inf. Process. Lett., 84(1):1–3, 2002.

[24] Ursula Molter and Ezequiel Rela. Furstenberg sets for a fractal set of directions. Proc. Amer. Math. Soc., 140:2753–2765, 2012.

[25] Andre Nies. Computability and Randomness. Oxford University Press, Inc., New York, NY, USA, 2009.

[26] Jan Reimann. Computability and fractal dimension. PhD thesis, Heidelberg University, 2004.

[27] Jan Reimann. Effectively closed sets of measures and randomness. Annals of Pure and Applied Logic, 156(1):170–182, 2008.

[28] Elias M. Stein and Rami Shakarchi. Real Analysis: Measure Theory, Integration, and Hilbert Spaces. Princeton Lectures in Analysis. Princeton University Press, 2005.
[29] Dennis Sullivan. Entropy, Hausdorff measures old and new, and limit sets of geometrically finite Kleinian groups. *Acta Mathematica*, 153(1):259–277, 1984.

[30] Claude Tricot. Two definitions of fractional dimension. *Mathematical Proceedings of the Cambridge Philosophical Society*, 91(1):57–74, 1982.