New Bounds for the Integer Carathéodory Rank

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Abstract. Given a rational pointed \(n\)-dimensional cone \(C\), we study the integer Carathéodory rank \(\text{CR}(C)\) and its asymptotic form \(\text{CR}^a(C)\), where we consider “most” integer vectors in the cone. The main result significantly improves the previously known upper bound for \(\text{CR}^a(C)\). We also study bounds on \(\text{CR}(C)\) in terms of \(\Delta\), the maximal absolute \(n \times n\) minor of the matrix given in an integral polyhedral representation of \(C\). If \(\Delta \in \{1, 2\}\), we show \(\text{CR}(C) = n\), and prove upper bounds for simplicial cones, improving the best known upper bound on \(\text{CR}(C)\) for \(\Delta \leq n\).

1 Introduction

A cone \(C\) in \(\mathbb{R}^n\) is \textit{rational} if there exists an integer \(m \times n\) matrix \(A\) such that

\[ C = \{ x \in \mathbb{R}^n : Ax \geq 0 \}, \quad (1) \]

where \(0\) is the zero vector and the inequality is componentwise. The cone \(C\) is \textit{pointed} if \(C \cap (-C) = \{0\}\), that is, the origin is the vertex of \(C\), or equivalently, \(A\) has full column rank. The \textit{dimension} of the cone \(C\) is the cardinality of a maximal set of linearly independent vectors in \(C\).

Given a finite set \(G = \{g_1, \ldots, g_t\} \subset \mathbb{Z}^n\), the \textit{semigroup} \(S\) with \textit{generating set} \(G\) is defined as

\[ S = \{ \lambda_1 g_1 + \cdots + \lambda_t g_t : \lambda_1, \ldots, \lambda_t \in \mathbb{Z}_{\geq 0} \}. \quad (2) \]

Let \(C \subset \mathbb{R}^n\) be a rational pointed \(n\)-dimensional cone. The integer points in the cone \(C\) form a semigroup \(S = C \cap \mathbb{Z}^n\). Due to a result of van der Corput \[26\] (see also \[21\] and \[23\]), the semigroup \(S\) has a uniquely determined inclusion–minimal finite generating set \(H = H(C)\). The set \(H\) has been traditionally referred to as the \textit{Hilbert basis} of \(C\). In the theory of mathematical optimisation, Hilbert bases are strongly related to \textit{Totally Dual Integral (TDI)-systems} \[24, \text{Chapter 22.3}\], and \textit{Graver bases} \[1, 17\].
A classical theorem by Carathéodory states that each point of the cone \( C \) is a non-negative combination of at most \( n \) vectors which lie on extreme rays of \( C \). Cook, Fonlupt, and Schrijver posed in [12] the following question, analogous to the one answered by Carathéodory’s theorem:

- What is the smallest \( k \) such that every integer point of the cone \( C \) can be expressed as a non-negative integer combination of at most \( k \) vectors in the Hilbert basis \( H \)?

We will refer to \( k \) as the integer Carathéodory rank of \( C \), and denote it by \( CR(C) \). To formally define this quantity, we consider for any element \( x \) of the semigroup \( C \cap \mathbb{Z}^n \) its representation length

\[
\sigma(x) = \min \{ l : x = \lambda_1 h_1 + \cdots + \lambda_l h_l, \lambda_i \in \mathbb{Z}_{\geq 0}, h_i \in H(C) \}.
\]

Then

\[
CR(C) = \max \{ \sigma(x) : x \in C \cap \mathbb{Z}^n \}.
\]

The paper [12] gave the upper bound \( CR(C) \leq 2n - 1 \) which was, subsequently, applied in the context of TDI-systems, integer rounding property of integer programs, independent sets of matroids, and coverings of perfect graphs. The current best known upper bound

\[
CR(C) \leq 2n - 2 \quad (3)
\]

was obtained by Sebő in [25]. We also remark that Bruns and Gubeladze [7] studied the maximum representation length for the elements of general semigroups \( S \) of the form (2). We refer the reader to [2, 3, 14] for known results in the general case.

Following the work of Bruns and Gubeladze [7], we say that a pointed rational \( n \)-dimensional cone \( C \subset \mathbb{R}^n \) satisfies the Integral Carathéodory Property (ICP) if \( CR(C) = n \). It was conjectured in [25] that the ICP holds for every \( n \)-dimensional cone \( C \). This conjecture was disproved by Bruns et al. in [8]. Specifically, it was shown in [8] that in every dimension \( n \geq 6 \) there exists an \( n \)-dimensional cone \( C \) with \( CR(C) \geq \lceil 7n/6 \rceil \).

To study the “typical” maximal representation length, Bruns and Gubeladze introduced in [7] the asymptotic integer Carathéodory rank \( CR^a(C) \) of the cone \( C \), which is defined as the smallest positive integer \( k \) such that the following limit exists and satisfies the equality

\[
\lim_{\delta \to \infty} \frac{|\{x \in S : \sigma(x) \leq k\} \cap [-\delta, \delta]^n|}{|S \cap [-\delta, \delta]^n|} = 1.
\]

That is “most” vectors in \( C \cap \mathbb{Z}^n \) can be represented by at most \( k \) Hilbert basis elements.

Clearly, \( CR^a(C) \leq CR(C) \). It was shown in [7] that

\[
CR^a(C) \leq 2n - 3 \quad (4)
\]

and that in every dimension \( n \geq 6 \) there exists an \( n \)-dimensional cone \( C \) with \( CR^a(C) > n \).
Known results on the integer Carathéodory rank lead to two interesting and long-standing open questions:

- **What are the optimal upper bounds for** \( \text{CR}(C) \text{ and CR}^a(C) \text{ in terms of } n? \) In the case of \( \text{CR}(C) \), Sebő's bound (3) remains the best known upper estimate for over three decades. Further, the work of Gubeladze [19] indicates that reducing (3) to a bound of the form \( \text{CR}(C) \leq (2 - \epsilon)n \) with \( \epsilon > 0 \) for all sufficiently large \( n \) is a challenging problem. In particular, it would disprove a conjecture ([19, Conjecture 2.1]) on the integer Carathéodory rank of normal polytopes.

In this paper, we study the above question in the case of the asymptotic integer Carathéodory rank. Theorem 1 reduces the bound (4) to \( \text{CR}^a(C) \leq \lfloor \frac{3n}{2} \rfloor \). On the other hand, Theorem 2 shows that in every dimension \( n \geq 6 \) there exists an \( n \)-dimensional cone \( C \) with \( \text{CR}^a(C) \geq \lfloor \frac{7n}{6} \rfloor \).

- **What cones have the integer Carathéodory property?** Cook, Fonlupt, and Schrijver [12] observed that the ICP holds for two-dimensional cones. Subsequently, Sebő [25] proved the ICP for cones of dimension three. On the other hand, due to the result of Bruns et al. [8], there exist cones that do not satisfy the ICP for every \( n \geq 6 \). Despite this, it remains an active line of research to classify which cones admit the ICP; see [13, 16] for some results concerning cones related to matroids. We continue this line of research by investigating the integer Carathéodory rank in terms of the parameter

\[
\Delta(A) = \max \{|\det B| : B \text{ is an } n \times n \text{ submatrix of } A\},
\]

where \( A \in \mathbb{Z}^{m \times n} \) has full column rank. We refer to \( A \) as \( \Delta \)-modular if \( \Delta(A) = \Delta \).

Recently, significant effort has been made to understand the computational complexity of integer programming problems defined by \( \Delta \)-modular matrices. Three key results in this area are given in [4, 15, 22]. This task motivated the study of polyhedral geometry depending on the parameter \( \Delta(A) \); see [5, 9, 10, 20] for an incomplete collection of results concerning the distance of optimal integral solutions of an integer linear program and optimal vertex solutions of the corresponding relaxation, the lattice width of lattice-free polyhedra, and the diameter of polyhedra. In some of the recent advancements, in particular [10, 20], Hilbert bases play a central role when proving novel upper bounds which solely depend on \( \Delta(A) \).

An intriguing special case is *simplicial cones*, that is cones, where \( A \in \mathbb{Z}^{n \times n} \) in (1) satisfies \( \det A \neq 0 \). Even for simplicial cones it is an open question whether they admit the ICP. An affirmative answer to this has, combined with some extra effort, the following strong implication: the integer vectors contained in the zonotope spanned by the primitive generators of a non-simplicial cone have the ICP. In addition to this, the study of simplicial cones and the ICP relates to various other concepts in mathematics such as simplices with the integer decomposition property which themselves are connected to weighted projective spaces; see for instance [6, 11].

In this paper, we consider the above question of which cones admit the ICP independently from the dimension of the cone. Theorem 3 shows that the ICP
holds in arbitrary dimension for cones with $\Delta(A) \in \{1, 2\}$. In Theorem 4 we further strengthen this for simplicial cones and obtain an improvement on the bound (3) if $\Delta(A) \leq n$.

In what follows, by $\text{int} X$ we denote the interior of a set $X$ and $\text{aff} X$ is the affine hull of $X$. We use the notation $[m]$ for the set $\{1, \ldots, m\}$. Given $A \in \mathbb{Z}^{m \times n}, I \subset [m]$, and $J \subset [n]$, we denote by $A_{I,J}$ the submatrix of $A$ with rows indexed by $I$ and columns indexed by $J$. If $J = [n]$, we write $A_I$, and similarly $A_{-J}$ when $I = [m]$. In the same manner, given a vector $x \in \mathbb{R}^n$ and a set $I \subset [n]$, we denote by $x_I \in \mathbb{R}^{|I|}$ the vector with coordinates indexed by $I$. The support of $x$ is defined as $\text{supp}(x) = \{i \in [n] : x_i \neq 0\}$. We denote by $\text{GL}(n, \mathbb{Z})$ the group of all $n \times n$ unimodular matrices, that is $A \in \mathbb{Z}^{n \times n}$ and $|\det A| = 1$. The standard unit vectors in $\mathbb{R}^n$ are denoted by $e_1, \ldots, e_n$.

2 Statement of results

Our main result strengthens the bound (4) obtained by Bruns and Gubeladze [7].

**Theorem 1.** Let $C$ be a rational pointed $n$-dimensional cone. Then $$\text{CR}^a(C) \leq \left\lfloor \frac{3}{2} n \right\rfloor.$$ The second result gives a new lower bound for the maximal value of the asymptotic integer Carathéodory rank of an $n$-dimensional cone.

**Theorem 2.** For every integer $n \geq 6$ there exists a rational pointed $n$-dimensional cone $C_n$ such that $$\text{CR}^a(C_n) \geq \left\lfloor \frac{7}{6} n \right\rfloor. \quad (5)$$

To state our parameterized results, we consider a rational pointed cone $C(A) = \{x \in \mathbb{R}^n : Ax \geq 0\}$ and estimate its integer Carathéodory rank in terms of the parameter $\Delta(A)$.

Firstly, we show that the ICP holds for the cone $C(A)$ if all $n \times n$ subdeterminants of $A$ are bounded by two.

**Theorem 3.** Let $A \in \mathbb{Z}^{m \times n}$ be a matrix of full column rank with $\Delta(A) \leq 2$. Then $\text{CR}(C(A)) = n$.

Note that a counterexample for Sebő’s conjecture obtained in [8] has a polyhedral representation with $\Delta(A) = 144$. Thus, the smallest value of $\Delta(A)$ for which the ICP fails has to be between 3 and 144.

Suppose now that $A \in \mathbb{Z}^{n \times n}$ is a nonsingular matrix. Then $C(A)$ is a simplicial cone and $\Delta(A) = |\det A|$. In this setting, we obtain an upper bound for $\text{CR}(C(A))$ which combines the parameter $\Delta(A)$ with the dimension $n$. This results in an improvement on the bound (3) for $\Delta(A) \leq n$.

**Theorem 4.** Let $A \in \mathbb{Z}^{n \times n}$ be a nonsingular matrix.
If \( 1 \leq \Delta(A) \leq 4 \), then \( \text{CR}(C(A)) = n \).

(ii) If \( \Delta(A) \geq 5 \), then \( \text{CR}(C(A)) \leq n + \Delta(A) - 3 \).

3 Proofs of Theorem 1 and 2

3.1 Proof of Theorem 1

To prove the upper bound for the asymptotic integer Carathéodory rank \( \text{CR}(a) \), it is sufficient to construct a set \( D \subset C \cap \mathbb{Z}^n \) that satisfies the following two properties.

(i) For any point \( b \in D \) we have \( \sigma(b) \leq 3n/2 \).

(ii) We have \( \lim_{\delta \to \infty} \frac{|D \cap [-\delta, \delta]^n|}{|C \cap \mathbb{Z}^n \cap [-\delta, \delta]^n|} = 1 \).

Let us take the Hilbert basis \( H(C) = \{ h_1, \ldots, h_t \} \) of the cone \( C \) and consider the matrix \( H \in \mathbb{Z}^{n \times t} \) with columns \( h_1, \ldots, h_t \). Then, given a \( b \in C \cap \mathbb{Z}^n \) we define \( Q(b) = \{ x \in \mathbb{R}^t_{\geq 0} : Hx = b \} \). Finding a representation of \( b \) with a small number of the Hilbert basis elements is equivalent to finding an integer vector in \( Q(b) \) with a small support size. Following the method introduced by Cook, Fonlupt, and Schrijver in [12], the idea is to construct a representation by decomposing an optimal vertex solution of the linear optimisation problem

\[
\max \left\{ x_1 + \cdots + x_t : x = (x_1, \ldots, x_t)^\top \in Q(b) \right\}.
\]

Since \( C \) is pointed and \( b \in C \), it is clear that (6) is feasible and bounded. In the asymptotic setting, we can discard integer points \( b \) that are “close” to certain boundary cases. This allows us to construct a – in a sense – complementary second decomposition. It turns out that at least one of the two decompositions must have small support size.

Let

\[ \Delta = \max \{ |\det H_{-I}| : I \subset [t], |I| = n \} \]

We define the set

\[
D = C \cap \mathbb{Z}^n \setminus \bigcup_{\tau \in \binom{[t]}{n-1}} \left\{ \sum_{i=1}^t \lambda_i h_i : 0 \leq \lambda_i \text{ for all } i \in \tau, 0 \leq \lambda_i < \Delta \text{ for all } i \in [t] \setminus \tau \right\}.
\]

The key property of \( D \) is that if \( b \in D \), then any vertex of \( Q(b) \) has large basic variables. In particular, this holds for an optimal vertex solution of (6).

To prove that \( D \) satisfies (ii), note that for each \( \tau \in \binom{[t]}{n-1} \) we can write

\[
\left\{ \sum_{i=1}^t \lambda_i h_i : 0 \leq \lambda_i \text{ for all } i \in \tau, 0 \leq \lambda_i < \Delta \text{ for all } i \in [t] \setminus \tau \right\} = P + C_\tau,
\]
Observe that $|C \cap \mathbb{Z}^n \cap [-\delta, \delta]^n| \in \Theta(\delta^n)$. Similarly, for any $\tau \in \binom{[n]}{k-1}$, it holds that $|(P + C_\tau) \cap \mathbb{Z}^n \cap [-\delta, \delta]^n| \in \Theta(\delta^k)$, where $k = \text{rank } H_\tau \leq |\tau| < n$. It holds

$$\frac{|D \cap [-\delta, \delta]^n|}{|C \cap \mathbb{Z}^n \cap [-\delta, \delta]^n|} \geq 1 - \frac{\sum_{\tau \in \binom{[n]}{k-1}} |(P + C_\tau) \cap \mathbb{Z}^n \cap [-\delta, \delta]^n|}{|C \cap \mathbb{Z}^n \cap [-\delta, \delta]^n|}.$$ 

The latter ratio tends to zero, as $\delta$ tends to infinity. Hence, $D$ satisfies (ii).

To show that $D$ satisfies (i), it is sufficient to prove that for any $b \in D$, there exists a $x \in Q(b) \cap \mathbb{Z}^t$ with $|\text{supp}(x)| \leq 3n/2$. Let $\lambda$ be an optimal vertex solution for (6). Hence, $\lambda$ has at most $n$ non-zero entries and, renumbering the coordinates, we may assume that $\lambda_{n+1} = \ldots = \lambda_t = 0$. Furthermore, the condition $b \in D$ implies that $\lambda_1, \ldots, \lambda_n \geq \Delta$.

Let $\mu = (\lambda_1 - \lfloor \lambda_1 \rfloor, \ldots, \lambda_n - \lfloor \lambda_n \rfloor)^\top$. Consider the vector $r = H_{\lfloor n \rfloor} \mu \in C \cap \mathbb{Z}^n$. We can write $r = \sum_{i=1}^t \beta_i h_i$, with $\beta_i \in \mathbb{Z}_{\geq 0}$. Hence,

$$b = \sum_{i=1}^n \lambda_i h_i = \sum_{i=1}^n \lfloor \lambda_i \rfloor h_i + \sum_{i=1}^n \mu_i h_i = \sum_{i=1}^n \lfloor \lambda_i \rfloor h_i + \sum_{i=1}^t \beta_i h_i.$$ 

By the optimality of $\lambda$, we have $\sum_{i=1}^t \beta_i \leq \sum_{i=1}^n \mu_i$.

If $\sum_{i=1}^t \mu_i \leq n/2$, then at most $\lfloor n/2 \rfloor$ of the numbers $\beta_i$ can be non-zero and the result follows. To settle the case $\sum_{i=1}^t \mu_i > n/2$, we consider the vector $\gamma = ([\lambda_1] - \lambda_1, \ldots, [\lambda_n] - \lambda_n)^\top$. We have $s = H_{\lfloor n \rfloor} \gamma \in C \cap \mathbb{Z}^n$ and, consequently, we can write $s = \sum_{i=1}^t \delta_i h_i$ with $\delta_i \in \mathbb{Z}_{\geq 0}$. Let $q = |\det H_{\lfloor n \rfloor}| \leq \Delta$. Then $q^\gamma$ is integral by Cramer’s rule. Recall that we have $b \in D$ and, in particular, $\lambda_i \geq \Delta$ for all $i \in [n]$. Thus, we obtain that

$$\eta_i = \lambda_i - (q-1)\gamma_i = (\lambda_i + \gamma_i) - q\gamma_i$$

are non-negative integers for all $i \in [n]$. Now, we can express $b$ as

$$b = \sum_{i=1}^n \lambda_i h_i = \sum_{i=1}^n \eta_i h_i + (q - 1) \sum_{i=1}^n \gamma_i h_i$$

$$= \sum_{i=1}^n \eta_i h_i + (q - 1) \sum_{i=1}^t \delta_i h_i.$$ 

Again, by the maximality of $\lambda$, we have that

$$\sum_{i=1}^t \delta_i \leq \sum_{i=1}^n \gamma_i < n/2.$$
and thus less than \( \frac{n}{2} \) of the numbers \( \delta_i \) can be non-zero. This completes the proof, since \( b \) is the non-negative integral combination of at most \( 3\frac{n}{2} \) Hilbert basis elements.

3.2 Proof of Theorem 2

Let \( C \) be a pointed rational \( n \)-dimensional cone in \( \mathbb{R}^n \). Theorem 6.1 in [7] implies that if \( CR^n(C) = n \), then \( CR(C) = n \). The counterexample to the integer Carathéodory conjecture, shown in [8], provides a 6-dimensional cone \( C_6 \) with \( CR(C_6) = 7 \). Hence, by Theorem 6.1, we get the lower bound \( CR^n(C_6) > 6 \).

Moreover, the inequality \( CR^n(C_6) \leq CR(C_6) \) implies \( CR^n(C_6) = 7 \).

Lemma 4.4 in [7] shows that \( CR^n(C \times C') = CR^n(C) + CR^n(C') \). Following the construction in [8] and setting \( C_n = \bigtimes_{i=1}^{\left\lfloor \frac{n}{6} \right\rfloor} C_6 \times C' \), where \( C' \) is any pointed, full-dimensional, rational cone in \( \mathbb{R}^n \mod 6 \), we obtain a cone that satisfies (5).

4 Proofs of Theorems 3 and 4

Throughout this section, we work with the polytope

\[
P_1(A) = \{ x \in \mathbb{R}^n : 0 \leq Ax \leq 1 \},
\]

where \( 1 \) denotes the all-ones vector. Note that \( P_1(A) \) is full-dimensional if and only if \( C(A) \) is full-dimensional. Further, \( P_1(A) \) is bounded as \( A \) has full column rank.

When proving Theorem 3 and Theorem 4, we employ the following strategy: Firstly, we argue that \( P_1(A) \cap \mathbb{Z}^n \setminus \{0\} \neq \emptyset \). Then, given \( z \in \text{int} C(A) \cap \mathbb{Z}^n \), this implies the existence of some Hilbert basis element \( h \in H(C(A)) \) such that the point \( z - \lambda h \) for some \( \lambda \in \mathbb{Z}_{>0} \) is contained in the boundary of \( C(A) \); see Lemma 5. Next, we work with our new integer vector \( z - \lambda h \) and iterate this procedure using Lemma 6 below. Hence, we use at every step exactly one Hilbert basis element and the dimension of the face of \( C(A) \) which contains the current integer vector in the relative interior decreases by at least one. Consequently, our strategy results in expressing \( z \) as an integer combination of at most \( n \) Hilbert basis elements.

We emphasize that the outlined strategy, when passing from \( z \) to \( z - \lambda h \), uses the observation that the Hilbert basis of the face \( F \) corresponding to \( z - \lambda h \) coincides with \( F \cap H(C(A)) \), that is \( H(F) = F \cap H(C(A)) \).

We begin by proving the first step, that is, the existence of a Hilbert basis element \( h \) from above. By doing so, we exploit a crucial property of Hilbert basis elements: given \( h \in H(C) \) and \( y_1, y_2 \in C \cap \mathbb{Z}^n \) such that \( h = y_1 + y_2 \), then either \( y_1 = 0 \) or \( y_2 = 0 \); see, e.g., [24, Chapter 16.4] for some details.

**Lemma 5.** Let \( A \in \mathbb{Z}^{m \times n} \) be a full column rank matrix with rows \( a_1^\top, \ldots, a_m^\top \) such that \( P_1(A) \cap \mathbb{Z}^n \setminus \{0\} \neq \emptyset \). Given \( z \in \text{int} C(A) \cap \mathbb{Z}^n \), there exists a Hilbert
basis element \( h \in H(C(A)) \) and \( \lambda \in \mathbb{Z}_{>0} \) such that \( z - \lambda h \in C(A) \) and \( a_i^\top (z - \lambda h) = 0 \) for some \( i \in [m] \).

**Proof.** Let \( h \in P_1(A) \cap \mathbb{Z}^n \setminus \{0\} \) be chosen such that \( |\text{supp}(Ah)| \) is minimal among all vectors in \( P_1(A) \cap \mathbb{Z}^n \setminus \{0\} \). We observe that

\[
\lambda = \min_{i \in \text{supp}(Ah)} a_i^\top z
\]

already yields the claim for \( h \) as \( Ah \in \{0, 1\}^m \). So it suffices to argue that \( h \) is a Hilbert basis element.

Let \( y_1, y_2 \in C(A) \cap \mathbb{Z}^n \) be such that \( h = y_1 + y_2 \). It is sufficient to show that one of the vectors \( y_1, y_2 \) is zero. Since \( Ah \in \{0, 1\}^m \), we have \( Ay_i \in \{0, 1\}^m \) for \( i = 1, 2 \) as well. However, this implies that \( \text{supp}(Ay_i) \subset \text{supp}(Ah) \) for \( i = 1, 2 \). The minimality of \( |\text{supp}(Ah)| \) implies that either \( y_1 = 0 \) or \( y_2 = 0 \). □

Lemma 5 enables us to argue inductively over the dimension \( n \). To make this precise, we establish in the next result a representation for lower-dimensional faces and their minors. Let \( A \in \mathbb{Z}^{m \times n} \) be a matrix with full column rank and let \( b \in \mathbb{Z}^m \). In what follows, we consider the polyhedron \( P(A, b) = \{ x \in \mathbb{R}^n : Ax \leq b \} \) and define

\[
\text{gcd}(A) = \text{gcd}(|\det(A_I)| : I \subset [m] \text{ with } |I| = n).
\]

For \( A \in \mathbb{Z}^{m \times n} \) with full row rank, we set \( \text{gcd}(A) = \text{gcd}(A^\top) \).

**Lemma 6.** Let \( A \in \mathbb{Z}^{m \times n} \) be a matrix with full column rank and \( b \in \mathbb{Z}^n \). Further, let \( F_I = P(A, b) \cap (v + \ker(A_I)) \) be a \((n-k)\)-dimensional face of \( P(A, b) \) with \( \text{aff} F_I \cap \mathbb{Z}^n \neq \emptyset \), where \( I \subset [m] \) with \( |I| = k \) and \( A_I, v = b_I \) hold. Then, there exists a unimodular transformation \( U \in \text{GL}(n, \mathbb{Z}) \) and orthogonal projection \( \pi : \mathbb{R}^n \rightarrow \mathbb{R}^{n-k} \) with the following properties:

(i) \( \pi (U \cdot F_I) \) is a \((n-k)\)-dimensional polyhedron that admits a representation of the form \( \pi (U \cdot F_I) = P(A, b) \) with an integer vector \( b \) and integer matrix \( A \) which is at most \( \left| \frac{\Delta(A)}{\text{gcd}(A_I)} \right| \)-modular.

(ii) There exists a one-to-one mapping between \( F_I \cap \mathbb{Z}^n \) and \( \pi (U \cdot F_I) \cap \mathbb{Z}^{n-k} \).

**Proof.** We assume without loss of generality that \( I = \{1, \ldots, k\} \). There exists a unimodular transformation \( U \in \text{GL}(n, \mathbb{Z}) \), such that

\[
A_{I^c}U^{-1} = (H, 0)
\]

for some nonsingular matrix \( H \in \mathbb{Z}^{k \times k} \) that can be obtained, for instance, by transforming \( A_{I^c} \) into Hermite normal form; see [24, Chapter 4] for more information on the Hermite normal form. Moreover, we have

\[
AU^{-1} = \begin{pmatrix} H & 0 \\ A & \bar{A} \end{pmatrix}
\]
for some $\tilde{A} \in \mathbb{Z}^{(m-k)\times (n-k)}$ with full column rank and $\tilde{A} \in \mathbb{Z}^{(m-k)\times k}$.

Let $\pi : \mathbb{R}^n \to \mathbb{R}^{n-k}$ denote the orthogonal projection onto the last $n-k$ coordinates and let $\tilde{z} \in \mathbb{R}^k$ be the unique solution of $Hx = b_I$. Then,
\[
\pi (U \cdot F) = \left\{ x \in \mathbb{R}^{n-k} : \tilde{A}x \leq b_{m\setminus I} - \tilde{A}\tilde{z} \right\}
\]
which is a $(n-k)$-dimensional polyhedron defined by the integral constraint matrix $\tilde{A}$. Let us show that $\tilde{A}$ is at most $\left\lfloor \frac{\Delta(A)}{\gcd(A_{II})} \right\rfloor$-modular. For that purpose, let $B$ be a $(n-k) \times (n-k)$ submatrix of $\tilde{A}$. We can extend the matrix to
\[
\begin{pmatrix}
H & 0 \\
\ast & B
\end{pmatrix}
\]
which is an $n \times n$ submatrix of $AU^{-1}$ with determinant $|\det B| |\det H| \leq \Delta(A)$. We have $|\det H| = \gcd(A_{II})$, which follows, e.g., from the Smith normal form; see for instance [24, Chapter 4.4] for a treatment of Smith normal forms. The latter equality and the integrality of $\tilde{A}$ imply that $\tilde{A}$ is at most $\left\lfloor \frac{\Delta(A)}{\gcd(A_{II})} \right\rfloor$-modular. Further, the right-hand side of the system in (7) defining $\pi(U \cdot F)$ is given by $b_{m\setminus I} - \tilde{A}\tilde{z}$. To settle property (i), it remains to prove the integrality of the right-hand side. We claim that $\tilde{z}$ has to be integral. The integrality follows then from the previous mentioned description of the right-hand side. Since $U$ is unimodular, $\text{aff}(U \cdot F) \cap \mathbb{Z}^n \neq \emptyset$ implies $\text{aff}(U \cdot F_I) \cap \mathbb{Z}^n \neq \emptyset$. Select $z \in \text{aff}(U \cdot F_I) \cap \mathbb{Z}^n$. So we have $Hz = b_I$. Since the solution of $Hx = b_I$ is unique, we have $\tilde{z} = z_I \in \mathbb{Z}^k$. Hence, property (i) follows by the discussion above.

For property (ii), observe that
\[
\text{aff}(U \cdot F) = z + \ker A_I, U^{-1} = z + \left\{ x \in \mathbb{R}^n : x_{[k]} = 0 \right\} = \left\{ x \in \mathbb{R}^n : x_{[k]} = \tilde{z} \right\}.
\]
So we obtain $y \in U \cdot F_I \cap \mathbb{Z}^n$ if and only if $\pi(y) \in \pi(U \cdot F_I) \cap \mathbb{Z}^{n-k}$, which settles property (ii). \hfill $\square$

### 4.1 Proof of Theorem 3

In this proof, we refer to $A$ as unimodular and bimodular if $\Delta(A) = 1$ and $\Delta(A) = 2$, respectively.

Suppose first that $\Delta(A) = 1$. We argue inductively on $n$. For $n = 1$, the cone is a ray and, thus, the statement immediately holds. Let $n > 1$ and $z \in C(A) \cap \mathbb{Z}^n$. If $z$ lies on the boundary of $C(A)$, we restrict to the face which contains $z$ in the relative interior and apply Lemma 6 with respect to that face. This results in a lower-dimensional cone with unimodular constraint matrix and the statement of the theorem follows by induction.

We may now assume $z \in \text{int} C(A) \cap \mathbb{Z}^n$, which implies that $C(A)$ is a full-dimensional cone. Since $P_1(A)$ is defined by a unimodular matrix with integral
right-hand side, every vertex of $P_1(A)$ is integral. There are at least two vertices as $n \geq 2$. Hence, the polytope $P_1(A)$ contains a non-zero integral vector. Using Lemmas 5 and 6, the statement of the theorem follows by induction.

Suppose now that $\Delta(A) = 2$. In this case, our proof crucially relies on an integer feasibility result due to Veselov and Chirkov:

**Theorem 7.** [27, Theorem 1] Let $A \in \mathbb{Z}^{m \times n}$ be bimodular and $b \in \mathbb{Z}^m$ such that $P(A, b)$ is full-dimensional. Then $P(A, b) \cap \mathbb{Z}^n \neq \emptyset$.

We argue again inductively over $n$. Similarly to the $\Delta(A) = 1$ case, we may assume that $n > 1$ and $z \in \text{int } C(A) \cap \mathbb{Z}^n$. Furthermore, we assume that every row of $A$ defines a facet of $C(A)$. If this is not the case, we remove rows of $A$ which do not correspond to a facet of $C(A)$. This operation does not increase $\Delta(A)$.

As before, our first goal is to show that $P_1(A)$ contains a non-zero integer vector. Let $F_{a} = \{x \in P_1(A) : a^T x = 1\}$ define a facet of $P_1(A)$ for some row $a$ of $A$. In what follows, we distinguish between the cases $\gcd(a) = 1$ and $\gcd(a) = 2$. Note that $\gcd(a) \geq 3$ is not possible as it violates the assumption that $A$ is bimodular.

Let $\gcd(a) = 1$. This implies that the affine hull of the facet $F_a$ contains integer vectors. Applying Lemma 6, we receive a full-dimensional polytope in $(n-1)$-dimensional ambient space which is bimodular. Hence, it contains an integer vector by Theorem 7. Therefore, the facet $F_a$ contains an integer vector by Lemma 6. Since the right-hand side equals one, this vector cannot be $0$.

Let $\gcd(a) = 2$. Since $n \geq 2$, there exists a different facet $F_{\tilde{a}} = \{x \in \mathbb{R}^n : \tilde{a}^T x = 1\}$. We have $\gcd(\tilde{a}) = 1$, otherwise we can choose a submatrix $A_I$, for $I \subseteq [m]$ with $|I| = n - 2$ such that $|\det(\tilde{a}, A_I)| \geq \gcd(a) \cdot \gcd(\tilde{a}) = 4$, a contradiction to the assumption that $A$ is bimodular. So we apply the previous argument to $F_{\tilde{a}}$ and obtain a non-zero integer vector in this facet.

Using Lemmas 5 and 6, the statement of the theorem follows by induction. □

### 4.2 Proof of Theorem 4

Recall that $A \in \mathbb{Z}^{n \times n}$ is a nonsingular matrix and, thus, $\Delta(A) = |\det A|$. In this manner, we write $|\det A|$ instead of $\Delta(A)$ throughout the proof. We will need the auxiliary result below. The proof is based on the theory of lattices; see [18] for an introduction to lattices.

**Lemma 8.** Suppose that $n \geq |\det A|$. Then the parallelepiped $P_1(A) = \{x \in \mathbb{R}^n : 0 \leq Ax \leq 1\}$ contains a non-zero integer vector.

**Proof.** Suppose that the matrix $A^{-1}$ has columns $w_1, \ldots, w_n$ and let $\Lambda = A^{-1} \mathbb{Z}^n$. Then $w_i \in \Lambda$ for each $i \in [n]$. Note that $A w_i = e_i$ and hence $w_1, w_1 + w_2, \ldots, w_1 + \cdots + w_n \in P_1(A) \cap \Lambda$. We analyse these sums with respect to the cosets of the finite abelian group $\Lambda/\mathbb{Z}^n$. Note that $|\Lambda/\mathbb{Z}^n| = |\det A|$ since $\det \Lambda = |\det A|^{-1}$. As $n \geq |\det A|$, either one of the sums is integral or
two sums, say $w_1 + \cdots + w_p$ and $w_1 + \cdots + w_q$ for $p < q$, are contained in the same coset of $\Lambda/\mathbb{Z}^n$ by the pigeonhole principle. This implies that $w_{p+1} + \cdots + w_q \in P_1(\mathbf{A}) \cap \mathbb{Z}^n$.

Suppose first that $1 \leq |\det \mathbf{A}| \leq 4$ and take any integer point $z$ in $C(\mathbf{A})$. By Lemma 6, we may assume that $C(\mathbf{A})$ is full-dimensional and that $z \in \text{int} C(\mathbf{A}) \cap \mathbb{Z}^n$. Since every cone of dimension at most three has the ICP [25, Theorem 2.2], we suppose that $n \geq 4$. Thus, we get $n \geq 4 \geq |\det \mathbf{A}|$. As $n \geq |\det \mathbf{A}|$, Lemma 8 gives $P_1(\mathbf{A}) \cap \mathbb{Z}^n \setminus \{0\} \neq \emptyset$. Applying Lemmas 5 and 6 we replace $z$ with an integer point in a lower-dimensional cone whose constraint matrix is at most $|\det \mathbf{A}|$-modular. We repeat this procedure until the dimension is at most $|\det \mathbf{A}| - 1$. For each iteration, we use exactly one Hilbert basis element and the number of iterations is at most $n - (|\det \mathbf{A}| - 1)$. Next, we apply again the result of [25] stating that the ICP holds for cones in dimension $3 \geq |\det \mathbf{A}| - 1$. Thus, we obtain an expression of $z$ as an integer combination of at most $n$ elements of $H(C(\mathbf{A}))$.

Suppose now that $|\det \mathbf{A}| \geq 5$ and take any integer point $z$ in $C(\mathbf{A})$. Observe that the case $n - 1 \leq |\det \mathbf{A}|$ follows from (3). Hence, we may assume $n \geq |\det \mathbf{A}|$. Therefore, using Lemma 8 and then Lemmas 5 and 6 as above, we can replace $z$ with an integer point in a lower-dimensional cone which is at most $|\det \mathbf{A}|$-modular. We repeat this procedure until the dimension is at most $|\det \mathbf{A}| - 1$. As $|\det \mathbf{A}| - 1 \geq 2$, we can apply Sebő's bound, (3), and get at most $2(|\det \mathbf{A}| - 1) - 2$ Hilbert basis elements in an integral combination. Together with our previous steps, which give us at most $n - (|\det \mathbf{A}| - 1)$ Hilbert basis elements in an integral combination, we obtain an expression of $z$ as a non-negative integer combination of at most

$$2(|\det \mathbf{A}| - 1) - 2 + (n - (|\det \mathbf{A}| - 1)) = n + |\det \mathbf{A}| - 3$$

elements of $H(C(\mathbf{A}))$. 

\textbf{Remark:} We highlight the limitation of our approach: Given an integer vector $z \in C(\mathbf{A})$, the strategy of searching for an element $h \in H(C(\mathbf{A}))$ such that for some integer $\lambda$ the vector $z - \lambda h$ reaches a lower-dimensional face of the cone $C(\mathbf{A})$ is limited to the case $\Delta(\mathbf{A}) \leq 2$. This fails for the well-understood case when $\Delta(\mathbf{A}) = 3$ and $n = 2$, which, in light of our previous discussion, could be considered the natural next step towards a possible extension of the method. An instance illustrating this deficiency is given by $C(\mathbf{A})$ with

$$\mathbf{A} = \begin{pmatrix} 1 & 0 \\ 2 & 3 \end{pmatrix}$$

and $z = (7, -3)^T \in C(\mathbf{A}) \cap \mathbb{Z}^n$. Then $\Delta(\mathbf{A}) = \det \mathbf{A} = 3$ and the Hilbert basis of $C(\mathbf{A})$ is given by the vectors $e_1, e_2, (2, -1)^T$, and $(3, -2)^T$. One can check that $P_1(\mathbf{A}) \cap \mathbb{Z}^n = \{0\}$ and no Hilbert basis element has the desired property. So in this comparably simple case the method of reducing to lower-dimensional faces of the cone fails.
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