ON THE EDIT DISTANCE FROM $K_{2,t}$-FREE GRAPHS

RYAN R. MARTIN AND TRACY MCKAY

Abstract. The edit distance between two graphs on the same vertex set is defined to be the size of the symmetric difference of their edge sets. The edit distance function of a hereditary property, $\mathcal{H}$, is a function of $p$, and measures, asymptotically, the furthest graph of edge density $p$ from $\mathcal{H}$ under this metric. In this paper, we address the hereditary property $\text{Forb}(K_{2,t})$, the property of having no induced copy of the complete bipartite graph with 2 vertices in one class and $t$ in the other. Employing an assortment of techniques and colored regularity graph constructions, we are able to determine the edit distance function over the entire domain $p \in [0,1]$ when $t = 3, 4$ and extend the interval over which the edit distance function for $\text{Forb}(K_{2,t})$ is known for all values of $t$, determining its maximum value for all odd $t$. We also prove that the function for odd $t$ has a nontrivial interval on which it achieves its maximum. These are the only known principal hereditary properties for which this occurs.

In the process of studying this class of functions, we encounter some surprising connections to extremal graph theory problems, such as strongly regular graphs and the problem of Zarankiewicz.

1. Introduction

The study of the edit distance in graphs initially appeared in a paper by Axenovich, Kézdy and the first author [2] and, independently, by Alon and Stav [1]. It has several potential applications, such as to biological consensus trees [2] and property testing problems in theoretical computer science [1]. More recently, interest has been shown in determining the value of the edit distance function, introduced in [3] by Balogh and the first author. Strategies for determining this function appear in [11], by Marchant and Thomason, as well as in [12].

Given a hereditary property (that is, a set of graphs closed under vertex deletion and isomorphism), what is the least number of edge additions or deletions sufficient to make any graph on $n$ vertices a member of the property? What is the behavior of this value as $n \to \infty$? In [2], the binary chromatic number of a graph $H$ is used as a means of bounding the maximum number of edge additions and deletions (edits) sufficient to ensure that every $n$-vertex graph has no induced copy of a single graph $H$. A hereditary property which consists of the graphs with no induced copy of $H$ is denoted $\text{Forb}(H)$ and is called a principal hereditary property. The methods in [2] give an asymptotically exact result in some cases, most notably when $H$ is self-complementary.

Let $G(n,p)$ denote the random graph on $n$ vertices with edge probability $p$. In [1], a version of Szemerédi’s regularity lemma is applied to show that, as $n \to \infty$, the number of edits necessary to make $G(n,p)$ a member of a given hereditary property approaches, with high probability, the maximum possible number over all $n$-vertex graphs within $o(n^2)$, so long as $p$ is chosen correctly with respect to the given hereditary property, $\mathcal{H}$. In fact, the maximum number of edits sufficient to change a density-$p$, $n$-vertex graph into a member of $\mathcal{H}$ is asymptotically the same as the expected number of edits required to put $G(n,p)$ into $\mathcal{H}$.

2010 Mathematics Subject Classification. Primary 05C35; Secondary 05C80.

Key words and phrases. edit distance, quadratic programming, strongly regular graphs.

The first author’s research partially supported by NSF grant DMS-0901008 and by an Iowa State University Faculty Professional Development grant.

The second author’s research partially supported from NSF grant DMS-0901008 and by an Iowa State University Mathematics Department Wolfe Research Fellowship.
The edit distance function, \( ed_{\mathcal{H}}(p) \), from [3], computes the limit of the maximum normalized edit distance of a density-\( p \), \( n \)-vertex graph from \( \mathcal{H} \) as \( n \to \infty \) for all probabilities \( p \). See Definition 2 below. Not surprisingly, the maximum value of this function occurs at the same \( p \) value described in [4].

Marchant and Thomason also explore the value of \( 1 - ed_{\mathcal{H}}(p) \) for various hereditary properties in [11], developing some insightful results for determining the value of the function in general. One discovery from [11] of particular interest is a relationship between the problem of determining the edit distance function for \( \text{Forb}(K_{3,3}) \) and constructions by Brown in [7] for \( K_{3,3} \)-free graphs, related to the Zarankiewicz problem.

More generally, edit distance has been discussed as a potential metric for graph limits, and also as a parameter for property testing techniques to which graph limits have been applied (see, for example, Borgs, et al. [4]).

In this paper, we explore what can be said about the edit distance function for the hereditary property \( \text{Forb}(K_{2,t}) \), the set of all graphs that do not contain a complete bipartite graph with cocliques of 2 and \( t \) vertices as an induced subgraph. In particular, we

- Compute the entire edit distance functions for the properties \( \text{Forb}(K_{2,3}) \) and \( \text{Forb}(K_{2,4}) \).
- Explore constructions that arise from strongly regular graphs and provide good upper bounds for \( ed_{\text{Forb}(K_{2,t})}(p) \). One in particular, derived from the 15-vertex generalized quadrangle \( GQ(2,2) \), defines \( ed_{\text{Forb}(K_{2,4})}(p) \) for \( p \in (1/5,1/3) \).
- Compute the edit distance function for the property \( \text{Forb}(K_{2,t}) \) for \( p \in \left[ \frac{2}{t+1}, 1 \right] \).
- Show that, for odd \( t \), \( ed_{\text{Forb}(K_{2,t})}(p) = \frac{1}{t+1} \) for \( p \in \left[ \frac{2t-1}{t(t+1)}, \frac{2t}{t+1} \right] \). These are the only known principal hereditary properties for which the maximum of the edit distance function is achieved on a nontrivial interval.
- Examine the relationship between constructions by Füredi [10] related to the Zarankiewicz problem and the trivial bound \( ed_{\text{Forb}(K_{2,t})}(p) \leq p(1-p) \), which achieves the value of the function for small values of \( p \) when \( t = 3, 4 \). When \( t \geq 9 \) the constructions by Füredi improve upon this bound for small values of \( p \).
- Look at constructions derived from powers of cycles that give a general upper bound for some values of \( p \) and \( t \).
- Derive a lower bound for \( ed_{\text{Forb}(K_{2,t})}(p) \) that is nontrivial and is achieved for some value of \( p \) if a specified strongly regular graph exists.
- Summarize the known bounds for \( ed_{\text{Forb}(K_{2,t})}(p) \) for \( 5 \leq t \leq 8 \).

Prior results and notation come primarily from [12], as well as previous work: [1], [2], [3], [11]; however, there are a number of other excellent resources on related topics. For a more extensive review of this literature, the reader may wish to consult Thomason [13]. We now introduce some important definitions and state our results more rigorously.

1.1. Definitions.

We begin by recalling the definitions of graph edit distance and the edit distance function.

**Definition 1** (Alon-Stav [1]; Axenovich-Kézdy-RM [2]). Let \( G \) and \( H \) be simple graphs on the same labeled vertex set, and let \( \mathcal{H} \) be a hereditary property, then

1. \( \text{dist}(G, H) = |E(G) \triangle E(H)| \) is the edit distance from \( G \) to \( H \),
2. \( \text{dist}(G, \mathcal{H}) = \min \{ \text{dist}(G, H) : H \in \mathcal{H} \} \) is the edit distance from \( G \) to \( \mathcal{H} \) and
3. \( \text{dist}(n, \mathcal{H}) = \max \{ \text{dist}(G, \mathcal{H}) : |G| = n \} \) is the maximum edit distance from the set of all \( n \)-vertex graphs to the hereditary property \( \mathcal{H} \).
Since $\mathcal{H}$ is by definition closed under isomorphism, vertex labels may be ignored when considering $\text{dist}(G, \mathcal{H})$. In fact, $\text{dist}(G, \mathcal{H})$ could be defined equivalently as the minimum number of edge changes necessary to make $G$ a member of $\mathcal{H}$.

The limit of the maximum edit distance from an $n$-vertex graph to a hereditary property $\mathcal{H}$ normalized by the total number of potential edges in an $n$-vertex graph,

$$d_\mathcal{H}^*= \lim_{n \to \infty} \text{dist}(n, \mathcal{H})/\binom{n}{2},$$

is demonstrated in [1] to exist and to be realized asymptotically with high probability by the random graph $G(n, p^*)$, where $p^* \in [0, 1]$ is a probability that depends on $\mathcal{H}$ and is not necessarily unique.

**Definition 2 ([3]).** The edit distance function of a hereditary property $\mathcal{H}$ is defined as follows:

$$ed_\mathcal{H}(p) = \lim_{n \to \infty} \max \{ \text{dist}(G, \mathcal{H}) : |V(G)| = n, |E(G)| = \binom{p(n)}{2} \}. $$

This function has also been denoted as $g_\mathcal{H}(p)$ in, for example, [3]. It was also proven in [3] that, if $E$ denotes the expectation, then

$$ed_\mathcal{H}(p) = \lim_{n \to \infty} E[\text{dist}(G(n, p), \mathcal{H})]/\binom{n}{2}. $$

The limits above were proven to exist in [3], and furthermore, $ed_\mathcal{H}(p)$ is both continuous and concave down. As a result, the edit distance function attains a maximum value that is equal to $d_\mathcal{H}^*$. The point, or interval, at which $d_\mathcal{H}^*$ is attained is denoted $p_\mathcal{H}^*$, and when it is evident from context, the subscript $\mathcal{H}$ may be omitted from both.

**Colored regularity graphs (CRGs)** are the building blocks for determining the edit distance function for specific hereditary properties. We leave the formal definition of CRGs as well as basic facts about them to Section 1.3 for the reader who is unfamiliar with these objects. A few of the new results given in Section 1.2 are structural in nature and so an understanding of CRGs is useful in order to put the results in full context.

### 1.2. New Results.

In this paper, we prove the following results for the hereditary properties $\text{Forb}(K_{2,3})$ and $\text{Forb}(K_{2,4})$. The case of $K_{2,2}$ is mentioned in Section 5.3 of [11].

**Theorem 3.** Let $\mathcal{H} = \text{Forb}(K_{2,3})$. Then $ed_\mathcal{H}(p) = \min\{p(1 - p), \frac{1 - p}{2} \}$ with $p_\mathcal{H}^* = \frac{1}{2}$ and $d_\mathcal{H}^* = \frac{1}{4}$.

**Theorem 4.** Let $\mathcal{H} = \text{Forb}(K_{2,4})$. Then $ed_\mathcal{H}(p) = \min\{p(1 - p), \frac{7p+1}{15}, \frac{1-2p}{3} \}$ with $p_\mathcal{H}^* = \frac{1}{3}$ and $d_\mathcal{H}^* = \frac{2}{5}$.

It should be noted that $p_\mathcal{H}^*$ and $d_\mathcal{H}^*$ from Theorem 3 could be found using alternative methods from previous literature as well. In fact, they are a direct result of Lemma 5.14 in [11], as is the value of $ed_\mathcal{H}(p)$ for $p \geq 1/2$ in both theorems. The $p_\mathcal{H}^*$ and $d_\mathcal{H}^*$ values in Theorem 3, however, are not so easily found. The techniques used to prove both theorems also have the potential to yield some results for $ed_\mathcal{H}(p)$ when $\mathcal{H} = \text{Forb}(K_{2,t})$ and $t \geq 5$, as discussed in Sections 6 and 7.

In [11], it is established that for $p \geq 1/2$ and $\mathcal{H} = \text{Forb}(K_{2,t})$, the edit distance function $ed_\mathcal{H}(p) = (1 - p)/(t - 1)$. We extend this result to hold true for $p \geq 2/(t + 1)$.

**Theorem 5.** Let $t \geq 4$, $p \geq 2/(t + 1)$ and $\mathcal{H} = \text{Forb}(K_{2,t})$, then $ed_\mathcal{H}(p) = (1 - p)/(t - 1)$.

This extension along with a new CRG construction results in the determination of $d_\mathcal{H}^*$, the maximum value of the edit distance function, for all odd $t$. Using the general lower bound in Theorem 5 below, we also demonstrate, via Theorem 7, that this maximum value occurs on a nondegenerate interval of values for $p$. That is, $p_\mathcal{H}^*$ is not a single value for all odd $t \geq 5$. 
Theorem 6. Let \( t \geq 3 \) and \( p < 1/2 \). If \( K \) is a black-vertex, \( p \)-core CRG with white and gray edges such that the gray edges have neither a \( K_{2,t} \) nor a \( B_{t-2} \) (as defined in Lemma 20), then

\[
g_K(p) \geq p - \frac{t-1}{4t-5} \left[ 3p - 2 + 2\sqrt{1-3p + (t+1)p^2} \right].
\]

Theorem 7. For odd \( t \geq 5 \) and \( \mathcal{H} = \text{Forb}(K_{2,t}) \),

\[
d^*_H = 1/(t+1) \quad \text{and} \quad p^*_H \geq \frac{2t-1}{t(t+1)}, \frac{2}{t+1}.
\]

For small \( p \) and \( t \) large enough, we demonstrate a similar result for Zarankiewicz type constructions by Füredi [10] to that discovered in [11] for \( \text{Forb}(K_{3,3}) \) and rejected for \( t = 3 \) and 4.

Theorem 8. For \( \mathcal{H} = \text{Forb}(K_{2,t}) \), the edit distance function \( e_{d\mathcal{H}}(p) \) ≤ \( \frac{t-1+p(2p^2-2q(t-1)-2t)}{2(q^2-1)} \) for any prime power \( q \) such that \( t-1 \) divides \( q-1 \).

Corollary 9. For \( t \geq 9 \), there exists a value \( q_0 \) so that if \( q > q_0 \), then \( \frac{t-1+p(2p^2-2q(t-1)-2t)}{2(q^2-1)} < p(1-p) \) for some values of \( p \), which approach 0 as \( q \) increases. That is, arbitrarily close to \( p = 0 \), there is some value for \( p \) such that \( e_{d\mathcal{H}}(p) < p(1-p) \).

A strongly regular graph construction provides the upper bound \( \frac{7t+1}{15} \) for \( e_{d\text{Forb}(K_{2,4})}(p) \). Such constructions continue to be relevant for larger \( t \) values.

Theorem 10. For any \((k,d,\lambda,\mu)\)-strongly regular graph, there exists a corresponding CRG, \( K \), such that

\[
f_K(p) = \frac{1}{k} + \frac{k-d-2}{k} p.
\]

If \( \lambda \leq t-3 \) and \( \mu \leq t-1 \), then \( K \) forbids \( K_{2,t} \) embedding, and when equality holds for both \( \lambda \) and \( \mu \),

\[
f_K(p) = \frac{t-1}{t-1 + d(d+1)} + \left( 1 - \frac{(d+2)(t-1)}{t-1 + d(d+1)} \right) p.
\]

There is a very close connection between the result in Theorem 6 and strongly regular graphs. If we take the expression in (2) and minimize it with respect to \( d \) (see expression 11), then we obtain the expression on the right-hand side of 11. In particular, we show that if a \((k,d,t-3,t-1)\)-strongly regular graph exists, then the corresponding CRG, \( K \), has \( f_K(p) \) tangent to the curve
Theorem 11. For $\mathcal{H} = \text{Forb}(K_{2,t})$,
\[
ed_{\mathcal{H}}(p) \leq \frac{3p + 1}{5 + t}.
\]

Observe that this bound is better than that of the trivial $\min\{p(1-p), (1-p)/(t-1)\}$ for $p \in \left(\frac{t+2-\sqrt{t^2-16}}{2t+10}, \frac{3}{2t+1}\right)$ when $t \geq 5$.

Below are known upper bounds for $5 \leq t \leq 8$. It should be noted that as our knowledge of existing strongly regular graphs increases, new upper bounds are also likely to be discovered.

Theorem 12. Let $\mathcal{H} = \text{Forb}(K_{2,t})$.

- If $t = 5$, then
  \[
ed_{\mathcal{H}}(p) \leq \min \left\{ p(1-p), \frac{1 + 75p}{96}, \frac{1 + 26p}{40}, \frac{1 + 5p}{6}, \frac{1 - p}{4} \right\}.
  \]

- If $t = 6$, then
  \[
ed_{\mathcal{H}}(p) \leq \min \left\{ p(1-p), \frac{1 + 63p}{85}, \frac{1 + 14p}{26}, \frac{1 + 7p}{17}, \frac{1 + 2p}{10}, \frac{1 - p}{5} \right\}.
  \]

- If $t = 7$, then
  \[
ed_{\mathcal{H}}(p) \leq \min \left\{ p(1-p), \frac{1 + 124p}{156}, \frac{1 + 76p}{100}, \frac{1 + 44p}{64}, \frac{1 + 31p}{49}, \frac{1 + 20p}{36}, \frac{1 + 5p}{16}, \frac{1 - p}{8}, \frac{1 - p}{6} \right\}.
  \]

- If $t = 8$, then
  \[
ed_{\mathcal{H}}(p) \leq \min \left\{ p(1-p), \frac{1 + 124p}{156}, \frac{1 + 95p}{125}, \frac{1 + 53p}{76}, \frac{1 + 20p}{36}, \frac{1 + 11p}{25}, \frac{1 + 5p}{16}, \frac{3p + 1}{13}, \frac{1 - p}{7} \right\}.
  \]

We compare these upper bounds to the lower bound in Theorem 6 via the figures in Appendix A.

1.3. CRGs and Background.

To help describe how $\ned_{\mathcal{H}}(p)$ may be calculated, some definitions from [1] are required.

Definition 13 (Alon-Stav [1]). A colored regularity graph (CRG), $K$, is a complete graph with vertices colored black or white, and with edges colored black, white or gray.

At times, it may be convenient to refer to the graph induced by edges of a particular color in a CRG, $K$. We shall refer to these graphs as the black, white and gray subgraphs of $K$. The investigation of edit distance in [11] and in [14] uses a different paradigm with an analogous structure called a type. Essentially, our black, white and gray are their red, blue and green, respectively.

Definition 14 (Alon-Stav [1]). A colored homomorphism from a (simple) graph $H$ to a colored regularity graph $K$ is a mapping $\phi : V(H) \mapsto V(K)$, which satisfies the following:

1. If $uv \in E(H)$, then either $\phi(u) = \phi(v) = m$ and $m$ is colored black, or $\phi(u) \neq \phi(v)$ and the edge $\phi(u)\phi(v)$ is colored black or gray.
2. If $uv \notin E(H)$, then either $\phi(u) = \phi(v) = m$ and $m$ is colored white, or $\phi(u) \neq \phi(v)$ and the edge $\phi(u)\phi(v)$ is colored white or gray.
Basically, a colored homomorphism is a map from a simple graph to a CRG so that black is only associated with adjacency, white is only associated with nonadjacency and gray is associated with adjacency, nonadjacency or both. We will refer to a colored homomorphism from a simple graph $H$ to a CRG $K$ as an embedding of $H$ in $K$, and we denote the set of all CRGs that only allow the embedding of simple graphs in a hereditary property $\mathcal{H}$ as $K(\mathcal{H})$ or merely $K$ when $\mathcal{H}$ is clear from the context. Since any hereditary property may be described by a set of forbidden induced subgraphs, an equivalent description of $K(\mathcal{H})$ is the set of all CRGs that do not permit the embedding of any of the forbidden induced subgraphs associated with $\mathcal{H}$. For instance, $K(\text{Forb}(K_{2,3}))$ is the set of all CRGs that do not admit $K_{2,3}$ embedding.

In order to calculate $ed_H(p)$, colored regularity graphs are used in $[3]$ in order to define the following functions:

\begin{align*}
(3) \quad f_K(p) &= \frac{1}{p^2} \left[ p(|\text{VW}(K)| + 2|\text{EW}(K)|) + (1 - p)(|\text{VB}(K)| + 2|\text{EB}(K)|) \right] \\
(4) \quad g_K(p) &= \min\{u^T M_K(p) u : u^T 1 = 1 \text{ and } u \geq 0\}.
\end{align*}

Here $K$ is a CRG with $k$ vertices. $\text{VW}(K)$, $\text{VB}(K)$, $\text{EW}(K)$ and $\text{EB}(K)$ represent the sets of white vertices, black vertices, white edges and black edges in $K$ respectively. $M_K$ is essentially a weighted adjacency matrix for $K$ with black vertices and edges receiving weight $1 - p$, white vertices and edges receiving weight $p$ and gray edges receiving weight $0$. From $[11]$, it is known that $ed_H(p) = \inf_{K \in K}\{f_K(p)\} = \inf_{K \in K}\{g_K(p)\}$. Moreover, Alon and Stav $[11]$ show that if $\chi_B$ is the binary chromatic number of $\mathcal{H}$, then $ed_H(1/2) = 1/(\chi_B - 1)$. Marchant and Thomason, demonstrate in $[11]$ that $ed_H(p) = \min_{K \in K}\{g_K(p)\}$. That is, given $p$ there exists at least one CRG, $K \in K$, such that $ed_H(p) = g_K(p)$.

If we say a CRG, $K$, is a sub-CRG of another CRG, $K'$, when $\text{VW}(K) \subseteq \text{VW}(K')$, $\text{VB}(K) \subseteq \text{VB}(K')$, $\text{EW}(K) \subseteq \text{EW}(K')$ and $\text{EB}(K) \subseteq \text{EB}(K')$, then it may be observed that $g_K(p) \geq g_{K'}(p)$. Furthermore, as was noted in $[11]$, if $g_K(p) = g_{K'}(p)$, then there is no need to consider both $K$ and $K'$ when attempting to determine $\min_{K \in K}\{g_K(p)\}$. Thus, a special subset of CRGs is defined as follows.

**Definition 15** (Marchant-Thomason $[11]$). A $p$-core CRG is a CRG $K'$ such that for no nontrivial sub-CRG $K$ of $K'$ is it the case that $g_K(p) = g_{K'}(p)$. In other words, if $K'$ is a $p$-core CRG, and $K$ is a nontrivial sub-CRG of $K'$, then $g_K(p) > g_{K'}(p)$.

It can be shown (see $[11]$) that a CRG, $K$, is $p$-core if and only if $g_K(p) = x^T M_K(p) x$ for a unique vector $x$ with positive entries summing to 1. Any CRG, $K$, that is not $p$-core contains at least one $p$-core sub-CRG $K'$ so that $g_{K'}(p) = g_K(p)$. Thus we could also say that

$$ed_H(p) = \min\{g_K(p) : K \in K \text{ and } K \text{ is } p\text{-core}\}.$$ 

That is, when looking for CRGs to determine $ed_H(p)$, the search may be limited to the important subset of CRGs, $p$-cores. This observation is especially helpful for determining lower bounds for the edit distance function.

To prove the upper bounds for the edit distance function in this paper, we show that for each $p$ there exists a CRG, $K \in K$, so that $f_K(p)$ gives the bound for the value of $ed_H(p)$. For the lower bounds, we employ so-called symmetrization, due to Sidorenko $[13]$, previously used for the computation of edit distance functions in $[12]$ and elsewhere such as Marchant and Thomason $[11]$. We also exploit features of the graphs $K_{2,1}$ and the concavity of the edit distance function to demonstrate that for no $p$-core CRG, $K \in K$, can $g_K(p)$ be less than the value in the theorem.

By the continuity of the edit distance function, if we know the value of the function on an open interval, then we also know the value on its closure. Hence, for convenience, most of our proofs will only address the value of the function on the interior of a given interval. We also note that, in $[12]$,
whenever the edit distance function of a hereditary property is computed, there is an attempt to
determine all of the $p$-core CRGs that achieve the value of the edit distance function. In this paper
we only concern ourselves with the value of the edit distance function itself and do not address the
issue of multiple defining constructions.

1.4. The Zarankiewicz problem and strongly regular graphs.

One reason for our interest in the edit distance function for Forb($K_{2,t}$) is its relation to the
Zarankiewicz problem. This problem addresses the question of how many edges a graph can have
before it must contain a $K_{s,t}$ subgraph for fixed $s$ and $t$. In an intriguing result from [11], a
construction from Brown [7] for $K_{3,3}$-free graphs is applied to construct an infinite set of new
CRGs that improve upon the previously known bounds for $ed_{Forb(K_{3,3})}(p)$ on certain intervals for
arbitrarily small $p$.

Marchant and Thomason [11] establish that it is sufficient to consider only $p$-core CRGs for
which the gray subgraph has neither a $K_3$ nor a $K_{3,3}$. Brown’s constructions are not $K_3$-free, but a
bipartite graph can be created from the construction that has no copy of $K_{3,3}$. Similarly, for $K_{2,t}$, it
is sufficient to have no gray $K_{2,t}$ or book $B_{t-2}$ to forbid $K_{2,t}$ embedding, where the graph $B_{t-2}$
is a “book” as defined in [8] and is defined precisely in Lemma 20.

Although the Brown constructions show that the edit distance function for Forb($K_{3,3}$) is strictly
less than $p(1-p)/(1+p)$ for sufficiently small $p$, known constructions for dense $K_{2,t}$-free graphs do
not play a role in the computation of the edit distance function for Forb($K_{2,3}$) or Forb($K_{2,4}$) in the
same way. However, the edit distance function for Forb($K_{2,4}$) is achieved over the interval $[1/5, 1/3]$
by a construction formed from a strongly regular graph, namely a generalized quadrangle, often
denoted $GQ(2,2)$, and similar results to those from the Brown constructions do reemerge when
t $\geq 9$ and Füredi’s constructions from [10] are considered.

1.5. Organization.

In Section 2 we discuss some results from [11] and [12] for the edit distance function and how
they may be applied to the problem of determining the function for Forb($K_{2,t}$). We then proceed
to some general results and observations in Section 3 that will be useful throughout the paper.
Sections 4 and 5 contain the proofs of our results for Forb($K_{2,3}$) and Forb($K_{2,4}$), respectively.
Section 6 addresses the proofs of Theorems 5, 6 and 7. In Section 7 we present several new CRG
constructions that yield upper bounds for $ed_{Forb(K_{2,t})}(p)$ in general. The remaining sections are
reserved for conclusions and acknowledgements.

2. Applications of past results to Forb($K_{2,t}$)

If a CRG is $p$-core, one can say some interesting things about its overall structure. From [11],
for instance, we have the following useful result.

**Theorem 16** (Marchant-Thomason [11]). If $K$ is a $p$-core CRG, then all edges of $K$ are gray except

- if $p < 1/2$, some edges joining two black vertices might be white or
- if $p > 1/2$, some edges joining two white vertices might be black.

The CRGs with all edges gray are useful in bounding the edit distance function, as we see in
[12].

**Definition 17.** Let $K(w,b)$ denote the CRG with $w$ white vertices, $b$ black vertices and only gray
edges. In particular:
(1) Let $K(1,1)$ be the CRG consisting of a white and black vertex joined by a gray edge.
(2) Let $K(0,t-1)$ be the CRG consisting of $t-1$ black vertices all joined by gray edges.

Theorem 18 settled the case of $K_{2,2}$, and Theorem 19 permits us to focus on black-vertex CRGs.

**Theorem 18** (Marchant-Thomason [11]). Let $\mathcal{H} =$ Forb($K_{2,2}$). Then $ed_\mathcal{H}(p) = g_{K(1,1)}(p) = p(1-p)$ with $p^*_\mathcal{H} = \frac{1}{2}$ and $d^*_\mathcal{H} = \frac{1}{2}$.

**Theorem 19** (Marchant-Thomason [11]). Let $\mathcal{H} = \text{Forb}(K_{2,t})$, $t > 2$. Then

(1) For $p > \frac{1}{2}$, $ed_\mathcal{H}(p) = g_{K(0,t-1)}(p) = \frac{1}{1-p}$ and

(2) For $p \leq \frac{1}{2}$, either

- $ed_\mathcal{H}(p) = \min\{g_{K(1,1)}(p), g_{K(0,t-1)}(p)\}$, or
- $ed_\mathcal{H}(p) = g_{K}(p) < \min\{g_{K(1,1)}(p), g_{K(0,t-1)}(p)\}$, where $K$ is a $p$-core CRG with only black vertices and, consequently, no black edges.

The following lemma is about the structure of the $p$-core CRGs described in the second part of Theorem 19. It was originally observed in Example 5.16 of [11]. The proof is a straightforward case analysis.

**Lemma 20** (Marchant-Thomason [11]). A CRG, $K$, on all black vertices with only white and gray edges forbids $K_{2,t}$ embedding if and only if its gray subgraph contains no $K_{2,t}$ or $B_{t-2}$ as a subgraph, where $B_{t-2}$ is a book as described in [8]. That is, the graph $B_{t-2}$ is defined to be the graph consisting of $t-2$ triangles that all share a single common edge.

As demonstrated in [11], for a $p$-core CRG, $K$, there is a unique vector $x$ so that $g_K(p) = x^T M_K(p)x$.

**Definition 21** (Marchant-Thomason [11]). For a $p$-core CRG $K$ with optimal weight vector $x$, the entry of $x$ corresponding to a vertex, $v \in V(K)$, is denoted by $x(v)$. This is the weight of $v$, and the function $x(v)$ is the optimal weight function.

With this in mind, we have two propositions from [12], which follow easily from [11].

**Proposition 22** ([12]). Let $K$ be a $p$-core CRG with all vertices black. Then for any $v \in V(K)$ and optimal weighting $x$, $d_G(v) = \frac{p - g_K(p)}{p} + \frac{1 - 2p}{p} x(v)$, where $d_G(v)$ is the sum of the weights of the vertices adjacent to $v$ via a gray edge.

**Proposition 23** ([12]). Let $K$ be a $p$-core CRG with all vertices black, then for $p \in [0,1/2]$ and optimal weighting $x$,

$$x(v) \leq \frac{g_K(p)}{1-p}, \forall v \in V(K).$$

Because of Theorem 19 we may restrict our attention to those CRGs, $K$, for which $g_K(p) \leq p(1-p)$. As a result, Proposition 22 gives the lower bound $d_G(v) \geq p + \frac{1 - 2p}{p} x(v)$. Meanwhile, Proposition 23 restricts the optimal weights of all vertices in $K$ to be no more than $p$. These two restrictions are useful when attempting to prove lower bounds for $ed_{\text{Forb}(K_{2,t})}(p)$.

3. Preliminary results and observations

We begin with some notation used throughout the paper.

**Definition 24.** Let $K$ be a black-vertex, $p$-core CRG with $g_K(p) \leq p(1-p)$ and optimal weight function $x$:

- $N_G(v) = \{y \in V(K) : vy \in E_G(K)\}$,
- $u_0$ is a fixed vertex in $K$ such that $x(u_0) \geq x(v)$, for all $v \in V(K)$, and $x = x(u_0)$ is its weight,
\[ U = N_G(u_0) \text{ and } |U| = \ell, \]
\[ u_1 \text{ is a fixed vertex with maximum weight in } U, \text{ and } x_1 = x(u_1), \]
\[ W \text{ is the set of all vertices in } K \text{ that are neither } u_0, \text{ nor contained in } U; \text{ or equivalently, } \]
\[ W \text{ is the set of all vertices in the white neighborhood of } u_0, \text{ and} \]
\[ x(S) = \sum_{y \in S} x(y) \text{ for some set } S \subseteq V(K). \]

Partitioning the vertices in a black-vertex, \( p \)-core CRG that forbids a \( K_{2,\ell} \) embedding into the three sets \( \{u_0\}, U \) and \( W \) as seen in Figure 3, illustrates some interesting features of its optimal weight function when the gray neighborhoods of these vertices are examined. One such feature is the upper bounds in Proposition 25 for \( x_1 \).

**Figure 3.** A partition of the vertices in a black-vertex, \( p \)-core CRG, \( K \). Dashed lines and gray background represent gray edges. White edges are omitted, as are edges within subsets.

**Proposition 25.** Let \( K \in [K(\text{Forb}(K_{2,3})) \cup K(\text{Forb}(K_{2,4}))] \) be a black-vertex, \( p \)-core CRG. If either \( p < 1/3 \) or both \( p < 1/2 \) and the gray sub-CRG of \( K \) is triangle-free, then

\[ x_1 \leq x \quad \text{and} \quad x_1 \leq p - x \]

where \( x = x(u_0) \) is the maximum weight of a vertex in \( K \), and \( x_1 = x(u_1) \) is the maximum weight of a vertex in that vertex’s gray neighborhood.

**Proof.** The inequality \( x_1 \leq x \) follows directly from definitions of \( x_1 \) and \( x \), since \( x \) is the greatest weight in \( K \). To justify the inequality \( x_1 \leq p - x \), we break the problem into two cases:

**Case 1:** \( u_0 \) and \( u_1 \) have no common gray neighbor.

Recall that \( u_1 \) is a vertex with maximum weight in the gray neighborhood of \( u_0 \), a vertex with maximum weight in all of \( K \), and assume that \( x + x_1 > p \). Then applying Proposition 22 and Theorem 19

\[ d_G(u_0) + d_G(u_1) \geq \left[ p + \frac{1 - 2p}{p} x \right] + \left[ p + \frac{1 - 2p}{p} x_1 \right] = 2p + \left( \frac{1 - 2p}{p} \right) (x + x_1) > 2p + (1 - 2p). \]

This is a contradiction because in Case 1, \( N_G(u_0) \cap N_G(u_1) = \emptyset \). Thus, \( d_G(u_0) + d_G(u_1) \leq 1 \), since the sum of the weights of the vertices in \( K \) must be 1.

This completes the proof of Proposition 25 for \( K \in K(\text{Forb}(K_{2,3})) \), since, in this case, no \( K \in K \) contains a gray triangle. So we may assume that \( K \in K(\text{Forb}(K_{2,4})). \)
Case 2: $u_0$ and $u_1$ have a common gray neighbor and $p < 1/3$.

In this case, $u_1$ has a single neighbor $u_2$ in $U$ because any more such neighbors would result in a gray book $B_2$. Furthermore, we note that in order to avoid a gray book $B_2$, the common neighborhood of $u_1$ and $u_2$ in $W$ must be empty. Consequently, $d_G(u_1) + d_G(u_2) \leq x(W) + 2x + x_1 + x(u_2)$.

Applying similar reasoning to that in Case 1,

$$d_G(u_0) + d_G(u_1) + d_G(u_2) \geq \left[ p + \frac{1 - 2p}{p}x \right] + \left[ p + \frac{1 - 2p}{p}x_1 \right] + \left[ p + \frac{1 - 2p}{p}x(u_2) \right].$$

So,

$$d_G(u_0) + (x(W) + 2x + x_1 + x(u_2)) \geq \left[ p + \frac{1 - 2p}{p}x \right] + \left[ p + \frac{1 - 2p}{p}x_1 \right] + \left[ p + \frac{1 - 2p}{p}x(u_2) \right].$$

$$x(U) + (x(W) + 2x + x_1 + x(u_2)) \geq 3p + \frac{1 - 2p}{p}(x + x_1 + x(u_2))$$

$$x(U) + x(W) + x \geq 3p + \frac{1 - 3p}{p}(x + x_1 + x(u_2))$$

$$1 \geq 3p + \frac{1 - 3p}{p}(x + x_1 + x(u_2)).$$

With $p < 1/3$ and $x + x_1 \geq p$, we have a contradiction. \hfill \Box

Applying the pigeon-hole principle, we also have the following lower bound for $\ell$:

**Fact 26.** In a CRG, if $u_0$ is a vertex with maximum weight, $x = x(u_0)$, the maximum weight in the gray neighborhood of $u_0$ is $x_1$, and the order of the gray neighborhood of $u_0$ is $\ell$, then $\ell \geq d_G(u_0)/x_1 \geq d_G(u_0)/x$.

While simple, when combined with Propositions 22 and 25 along with the observation that $x(u_0) + x(U) + x(W) = 1$, this fact forces a delicate balance between the weights of the vertex $u_0$, the vertices in $U$, and the vertices in $W$.

4. Proof of Theorem 3

In this section, we establish the value of $ed_{Forb(K_{2,3})}(p)$ for $p \in (0, 1/2)$, determining the entire function via continuity and Theorem 19 which gives that $ed_{Forb(K_{2,3})}(p) = (1 - p)/2$ for $p \in [1/2, 1]$.

For the following discussion, we assume that $K$ is a $p$-core CRG on all black vertices into which $K_{2,3}$ may not be embedded and that $g_K(p) \leq p(1 - p)$. The following lemma yields a useful restriction of the order of $U$.

**Lemma 27.** Let $K$ be a black-vertex, $p$-core CRG with $p \in (0, 1/2)$, no gray triangles, no gray $K_{2,3}$ and $g_K(p) \leq p(1 - p)$. If $u_0$ is a vertex of maximum weight, $x$, in $K$, and $\ell = |N_G(u_0)|$, then

$$\ell \leq \frac{2(1 - x) - \frac{1}{p}d_G(u_0)}{p - x}.$$ 

**Proof.** Let $u_1, \ldots, u_\ell$ be an enumeration of the vertices in $U$, the gray neighborhood of $u_0$. Observe that $K$ cannot contain a $K_3$ with all gray edges, and so $U$ contains no gray edges. Therefore, with the exception of $u_0$, the entire gray neighborhood of each $u_i$ is contained in $W$. Furthermore, if any three vertices in $U$ had a common gray neighbor in $W$, then $K$ would contain a gray $K_{2,3}$.
That is, each vertex in $W$ is adjacent to at most 2 vertices in $U$ via a gray edge. Applying these observations,

$$\sum_{i=1}^{\ell} (d_G(u_i) - x) \leq 2x(W).$$

Using Proposition 22 and the assumption that $\frac{p - g_K(p)}{p} \geq p$,

$$\sum_{i=1}^{\ell} \left( p - x + \frac{1 - 2p}{p}x(u_i) \right) \leq 2x(W).$$

The fact that $x(W) = 1 - x - d_G(u_0)$, gives

$$\ell(p - x) + \frac{1 - 2p}{p}d_G(u_0) \leq 2(1 - x - d_G(u_0))$$

$$\ell(p - x) \leq 2 - 2x - \frac{1}{p}d_G(u_0)$$

$$\ell \leq \frac{2(1 - x) - \frac{1}{p}d_G(u_0)}{p - x}.$$

□

The following technical lemma is an important tool in the proof of the Theorem.

**Lemma 28.** Let $K$ be a black-vertex, $p$-core CRG for $p \in (0, 1/2)$ with no gray triangles, no gray $K_{2,3}$ and $g_K(p) \leq p(1 - p)$. If $x$ and $x_1$ are defined as in Proposition 25, then

$$\left[ p + \frac{1 - 2p}{p}x \right] \left[ \frac{1}{x_1} + \frac{1}{p(p - x)} \right] \leq \frac{2(1 - x)}{p - x}.$$

**Proof.** By Fact 26 $\ell \geq \frac{d_G(u_0)}{x_1}$, and by Lemma 27 $\ell \leq \frac{2(1 - x) - \frac{1}{p}d_G(u_0)}{p - x}$. Therefore,

$$\frac{d_G(u_0)}{x_1} \leq \frac{2(1 - x) - \frac{1}{p}d_G(u_0)}{p - x}.$$

After combining the $d_G(u_0)$ terms we get,

$$d_G(u_0) \left[ \frac{1}{x_1} + \frac{1}{p(p - x)} \right] \leq \frac{2(1 - x)}{p - x},$$

and then applying Proposition 22

$$\left[ p + \frac{1 - 2p}{p}x \right] \left[ \frac{1}{x_1} + \frac{1}{p(p - x)} \right] \leq \frac{2(1 - x)}{p - x}.$$

□

We now turn to the proof of the main theorem for this section.

**Proof of Theorem 3.** Let $p \in (0, 1/2)$, and $K$ be a black-vertex, $p$-core CRG with $g_K(p) < p(1 - p)$ and no gray triangle (i.e., the book $B_1$) or gray $K_{2,3}$.

With the above assumptions, we will show that there is no possible value for $x$, the value of the largest vertex-weight. To do so, we break the problem into 2 cases: $x \geq \frac{p}{2}$ and $x < \frac{p}{2}$.

**Case 1:** $x \geq p/2$. 

...
We start with the inequality from Lemma 28,
\[
\left[ p + \frac{1-2p}{p} x \right] \left[ \frac{1}{x} + \frac{1}{p(p-x)} \right] \leq \frac{2(1-x)}{p-x},
\]
and apply the bound \(x_1 \leq p-x\) from Proposition 25 to get
\[
\left[ p + \frac{1-2p}{p} x \right] \left[ \frac{1}{p-x} + \frac{1}{p(p-x)} \right] \leq \frac{2(1-x)}{p-x}.
\]

From Proposition 23, \(p-x>0\), and so
\[
\left[ p + \frac{1-2p}{p} x \right] \left[ \frac{1}{x} + \frac{1}{p^2} \right] \left[ \frac{1-p}{p^2} \right] \leq 1 - p
\]
\[
x \leq p^2,
\]
a contradiction, since \(\frac{p}{2} > p^2\) for \(p \in (0, 1/2)\).

**Case 2**: \(x < p/2\).

We again apply Lemma 28, only now we employ the trivial bound \(x_1 \leq x\) from Proposition 25.
\[
\left[ p + \frac{1-2p}{p} x \right] \left[ \frac{1}{x} + \frac{1}{p(p-x)} \right] \leq \frac{2(1-x)}{p-x}
\]
\[
\left[ p + \frac{1-2p}{p} x \right] \left[ p(p-x) + x \right] \leq 2px(1-x)
\]
\[
(4p^2-3p+1)x^2 - (3p^3)x + p^4 \leq 0.
\]

Observe that \(4p^2-3p+1\) is always positive, and therefore the parabola \((4p^2-3p+1)x^2 - (3p^3)x + p^4\), in the variable \(x\), is concave up, so the range of \(x\) values for which this inequality is satisfied is \(x \in [x', x'']\) where
\[
x' = \frac{3p^3 - \sqrt{-4p^4 + 12p^5 - 7p^6}}{2(1 - 3p + 4p^2)} \quad \text{and} \quad x'' = \frac{3p^3 + \sqrt{-4p^4 + 12p^5 - 7p^6}}{2(1 - 3p + 4p^2)}.
\]

If \(p < (6 - 2\sqrt{2})/7\), then neither \(x'\) nor \(x''\) is real, and so the inequality is never satisfied. For \(p \in \left[\frac{6-2\sqrt{2}}{7}, \frac{1}{2}\right]\), routine calculations show that \(\frac{p}{2} < x'\), a contradiction to the assumption that \(x < \frac{p}{2}\).

Hence, there is no possible value for \(x\) if \(ed\binom{K_{2,3}}{p} < p(1-p)\), so the proof is complete. \(\Box\)

5. **Proof of Theorem 4**

This section addresses the case of \(ed\binom{K_{2,4}}{p}\).
5.1. Upper bounds.

Recall that from Theorem 19 we already know that $ed_{\text{Forb}(K_{2,4})}(p) \leq \min\{p(1-p), \frac{1+p}{3}\}$. For the remaining upper bound, we turn to strongly regular graphs:

**Definition 29.** A $(k, d, \lambda, \mu)$-**strongly regular graph** is a graph on $k$ vertices such that each vertex has degree $d$, each pair of adjacent vertices has exactly $\lambda$ common neighbors, and each pair of nonadjacent vertices has exactly $\mu$ common neighbors.

**Lemma 30.** Let $H = \text{Forb}(K_{2,t})$. If there exists a $(k, d, \lambda, \mu)$-strongly regular graph with $\lambda \leq t - 3$ and $\mu \leq t - 1$, then

$$ed_{H}(p) \leq \frac{1}{k} + \frac{k - d - 2}{k} p.$$  

**Proof.** Let $G$ be the aforementioned strongly regular graph. We construct a CRG, $K$, on $k$ black vertices with gray edges in $K$ corresponding to adjacent vertices in $G$ and white edges in $K$ corresponding to nonadjacent vertices in $G$.

No pair of adjacent vertices has $t - 2 > \lambda$ common neighbors, so there is no book $B_{t-2}$ in the gray subgraph, and no pair of vertices has $t > \mu, \lambda$ common neighbors, so there is no $K_{2,t}$ in the gray subgraph. Thus, by Lemma 20, $K_{2,t} \not\rightarrow K$. Furthermore,

$$f_{K}(p) = \frac{1}{k} \left[(1-p)k + 2p \left(\left(\frac{k}{2}\right) - \frac{dk}{2}\right)\right] = \frac{1}{k} + \frac{k - d - 2}{k} p.$$  

□

In fact, there is a $(15, 6, 1, 3)$-strongly regular graph [6]. It is a so-called “generalized quadrangle,” $\text{GQ}(2, 2)$. As a result,

$$ed_{\text{Forb}(K_{2,4})}(p) \leq \min\left\{p(1-p), \frac{1 + 7p}{15}, \frac{1 - p}{3}\right\}.$$  

A list of known strongly regular graphs and their parameters has been compiled by Andries Brouwer [6]. Their implications for $ed_{\text{Forb}(K_{2,4})}(p)$ in general are explored further in Section 7.1.

5.2. Lower bounds.

Because the edit distance function is both continuous and concave down, it is sufficient to verify that $ed_{\text{Forb}(K_{2,4})}(p) \geq p(1-p)$ for $p \in (0, 1/5)$ and that $ed_{\text{Forb}(K_{2,4})}(p) \geq (1-p)/3$ for $p \in (1/3, 1/2)$. This is because the line determined by the bound $\frac{1+7p}{15}$ passes through the points $(1/5, 4/25)$ and $(1/3, 2/9)$. Furthermore, by Theorem 19, we need only consider CRGs that have black vertices and white and gray edges.

Lemmas 31 and 34 settle the cases where $p \in (1/3, 1/2)$ and where $p \in (0, 1/5)$, respectively.

**Lemma 31.** Let $p \in (1/3, 1/2)$. If $K$ is a black-vertex, $p$-core CRG that does not contain a gray book $B_{2}$ or a gray $K_{2,4}$, then $g_{K}(p) \geq \frac{1-p}{3}$, with equality occurring only if $K$ is a gray triangle (i.e., $K \approx K(0,3)$).

**Proof.** See Appendix E

□

Before we prove Lemma 34 there are two propositions that are necessary and used in several cases.
Proposition 32. Let $p \in (0, 1/2)$, and let $K$ be a black-vertex, $p$-core CRG with no gray book $B_2$ and no gray $K_{2,4}$. If $g = g_K(p)$, $U = N_G(u_0)$, $\ell = |U|$ and $U_1 \subseteq U$ is the set of vertices in $U$ that are incident to a gray edge in $U$, then

$$\ell \left( \frac{p-g}{p} - x \right) \leq 3 - 3x - \frac{1+p}{p} x(U) + x(U_1) \leq 3 - 3x - \frac{1}{p} x(U).$$

Proof. See Appendix B. \hfill \square

Proposition 33. Let $p \in (0, 1/2)$, and let $K$ be a black-vertex, $p$-core CRG with no gray book $B_2$ and no gray $K_{2,4}$. If $g_K(p) \leq p(1-p)$, then both

$$p \geq \frac{9-4\sqrt{3}}{11} \quad \text{and} \quad x \geq \frac{p^2}{2(1-3p+5p^2)} \left[ 1 + 3p - \sqrt{3+18p-11p^2} \right] \geq \frac{1}{25}.$$ 

Proof. We begin with Proposition 32 and then use $\ell \geq x(U)/x$ from Fact 26

$$\ell \left( \frac{p-g}{p} - x \right) \leq 3 - 3x - \frac{1}{p} x(U),$$

$$\frac{x(U)}{x} \left( \frac{p-g}{p} - x \right) \leq 3 - 3x - \frac{1}{p} x(U),$$

$$x(U) \left( \frac{p-g}{px} - 1 + \frac{1}{p} \right) \leq 3 - 3x$$

(5)

$$\left[ \frac{p-g}{p} + \frac{1-2p}{p} x \right] \left[ \frac{p-g}{p} + \frac{1-p}{p} x \right] \leq 3x - 3x^2.$$

Recall that $(p-g)/p \geq p$ because $g \leq p(1-p)$, so

$$\left[ p + \frac{1-2p}{p} x \right] \left[ p + \frac{1-p}{p} x \right] \leq 3x - 3x^2$$

$$p^2 - (1+3p)x + \frac{1-3p+5p^2}{p^2}x^2 \leq 0.$$

The quadratic formula gives that not only must the discriminant be nonnegative (requiring $p \geq (9-4\sqrt{3})/11$), but also

$$x \geq \frac{p^2}{2(1-3p+5p^2)} \left[ 1 + 3p - \sqrt{3+18p-11p^2} \right].$$

Some routine but tedious calculations demonstrate that, for $p \in [(9-4\sqrt{3})/11, 1/2)$, this expression is at least 1/25, achieving that value uniquely at $p = 1/5$. \hfill \square

Lemma 34. Let $p \in (0, 1/5)$. If $K$ is a black-vertex, $p$-core CRG that does not contain a gray book $B_2$ or a gray $K_{2,4}$, then $g_K(p) > p(1-p)$.

Proof. We assume that $g_K(p) \leq p(1-p)$.

Case 1: $\ell \geq 8$.

According to Proposition 32

$$8 \left( \frac{p-g}{p} - x \right) \leq \ell \left( \frac{p-g}{p} - x \right) \leq 3 - 3x - \frac{1}{p} \left( \frac{p-g}{p} + \frac{1-2p}{p} x \right),$$

$$(1-2p-5p^2)x \leq 3p^2 - (p-g)(1+8p),$$
and since \( x \geq 1/25 \) and \( p - g \geq p^2 \),
\[
\frac{1 - 2p - 5p^2}{25} \leq 3p^2 - p^2(1 + 8p)
\]
\[
(1 - 5p^2)(1 + 8p) \leq 0,
\]
a contradiction. So, \( \ell < 8 \).

**Case 2:** \( \ell \leq 7 \) and \( x < p^2/(9p - 1) \).

Using Fact 26 and then Proposition 22
\[
7 \geq \ell \geq \frac{x(U)}{x} \geq \frac{p}{x} + \frac{1 - 2p}{p} > \frac{9p - 1}{p} + \frac{1 - 2p}{p} = 7,
\]
a contradiction.

**Case 3:** \( \ell \leq 7 \) and \( p^2/(9p - 1) \leq x \leq p/3 \).

First we bound \( \ell \):
\[
\ell \geq \frac{x(U)}{x} \geq \frac{p}{x} + \frac{1 - 2p}{p} \geq 3 + \frac{1}{p} - 2 > 6.
\]
So, \( \ell = 7 \). Since \( \ell \) is odd, \( x(U_1) \leq 6x \). By Proposition 32
\[
\ell \left( \frac{p - g}{p} - x \right) \leq 3 - 3x - \frac{1}{p}x(U) + x(U_1),
\]
and applying Proposition 22
\[
7 \left( \frac{p - g}{p} - x \right) \leq 3 - 3x - \frac{1}{p} \left[ \frac{p - g}{p} + \frac{1 - 2p}{p} x \right] + 6x
\]
\[
\frac{1 - p - 12p^2}{p^2} x \leq 3 - \frac{1 + 8p}{p} \cdot \frac{p - g}{p}
\]
\[
\frac{1 - p - 12p^2}{p^2} \left[ \frac{p^2}{9p - 1} \right] \leq 3 - \frac{1 + 8p}{p} \cdot p
\]
\[
\frac{1 - p - 12p^2}{p^2} \left[ \frac{9p - 1}{p - 1} \right] \leq 2(1 - 4p)
\]
\[
\frac{1 + 3p}{9p - 1} \leq 2,
\]
which implies \( p \geq 1/5 \), a contradiction.

**Case 4:** \( \ell \leq 7 \) and \( x > p/3 \).

Now we compute a stronger bound on \( U_1 \). Let \( u_1 \) and \( u_2 \) be vertices in \( U \) that are adjacent via a gray edge, and let their weights be \( x_1 \) and \( x_2 \), respectively. Then
\[
x + x(U) + (d_G(u_1) - x - x_2) + (d_G(u_2) - x - x_1) \leq 1
\]
because \( u_1 \) and \( u_2 \) have no common gray neighbor other than \( u_0 \) and because they can have no additional gray neighbor in \( U \). Applying Proposition 22
\[
x + \frac{p - g}{p} + \frac{1 - 2p}{p} x + 2 \frac{p - g}{p} - 2x + \frac{1 - 3p}{p} (x_1 + x_2) \leq 1
\]
\[
\frac{1 - 3p}{p} (x_1 + x_2) \leq \frac{3g - 2p}{p} - \frac{1 - 3p}{p} x,
\]
and since $p(1 - p) \geq g$,

$$x_1 + x_2 \leq p - x.$$  

We can bound the number of vertices in $U - U_1$ by using the fact that $(\ell - \ell_1)x \geq x(U) - x(U_1)$. Returning to Proposition 32,

$$\ell \left( \frac{p - g}{p} - x \right) \leq 3 - 3x - \frac{1 + p}{p}x(U) + x(U_1)$$

$$= \left[ \ell_1 + \frac{1}{x}x(U) - \frac{1}{x}x(U_1) \right] \left( \frac{p - g}{p} - x \right) \leq 3 - 3x - \frac{1 + p}{p}x(U) + x(U_1)$$

$$x(U) \left( \frac{p - g}{px} - 1 + \frac{1 + p}{p} \right) - 3 + 3x \leq x(U_1) \left( 1 + \frac{p - g}{px} - 1 \right) - \ell_1 \left( \frac{p - g}{p} - x \right).$$

If $\ell_1 = |U_1|$, then $x(U_1) \leq (\ell_1/2)(p - x)$. Of course, $x(U)$ is lower-bounded by Proposition 22

$$\left[ \frac{p - g}{p} + \frac{1 - 2p}{p}x \right] \left( \frac{p - g}{px} + \frac{1}{p} \right) - 3 + 3x \leq \frac{\ell_1}{2}(p - x) \left( \frac{p - g}{px} \right) - \ell_1 \left( \frac{p - g}{p} - x \right)$$

$$\left[ \frac{p - g}{p} + \frac{1 - 2p}{p}x \right] \left( \frac{p - g}{px} + \frac{1}{p} \right) - 3 + 3x \leq \ell_1 \left[ x - \frac{p - g}{p} \cdot \frac{3x - p}{2x} \right]$$

$$\left[ \frac{p - g}{p} + \frac{1 - 2p}{p}x \right] \left( \frac{p - g}{px} + \frac{1}{p} \right) - 3 + 3x \leq \ell_1 \left[ x - \frac{p(3x - p)}{2x} \right]$$

$$p^2 - (1 + 2p)x + \frac{1 - 2p + 3p^2}{p^2}x^2 \leq \ell_1 \frac{(p - x)(p - 2x)}{2}.$$  

Now, we bound $\ell_1$, depending on the sign of $p - 2x$, requiring two more cases.

**Case 4a**: $\ell \leq 7$ and $x > p/3$ and $p - 2x \geq 0$.

Here we use the bound $\ell_1 \leq 6$:

$$p^2 - (1 + 2p)x + \frac{1 - 2p + 3p^2}{p^2}x^2 \leq 3(p - x)(p - 2x)$$

$$-2p^2 + (7p - 1)x + \frac{1 - 2p - 3p^2}{p^2}x^2 \leq 0.$$  

By Proposition 23, we may restrict our attention to $p \geq (9 - 4\sqrt{3})/11 > 1/7$ and so we may substitute the smallest possible value for $x$, which still maintains the inequality.

$$-2p^2 + (7p - 1)\left( \frac{p}{3} \right) + \frac{1 - 2p - 3p^2}{p^2} \left( \frac{p}{3} \right)^2 < 0$$

$$-18p^2 + 3(7p - 1)p + (1 - 2p - 3p^2) < 0$$

$$1 - 5p < 0,$$

a contradiction.

**Case 4b**: $\ell \leq 7$ and $x > p/3$ and $p - 2x < 0$.  


Here we use the bound $\ell_1 \geq 0$ and then replace $x$ with $\frac{p^2(1+2p)}{2-4p+6p^2}$, the value that minimizes the left-hand side:

$$p^2 - (1+2p)x + \frac{1-2p+3p^2}{p^2} x^2 \leq 0$$
$$p^2 - \frac{(1+2p)^2p^2}{4(1-2p+3p^2)} \leq 0$$
$$p^2(3-12p+8p^2) \leq 0.$$

This, too, is a contradiction for $p \in (0,1/5)$, completing the proof of Lemma 34. □

6. Proofs of Theorems 5, 6 and 7

In this section we extend the generally known interval for $ed_{\text{Forb}(K_{2,t})}(p)$ from $p \in [1/2,1]$ to $p \in \left[\frac{2}{t+1},1\right]$. With a new CRG construction, this extension is sufficient to determine $d^*_p$ and a subset of $p^*_t$ for odd $t$. Subsection 6.1 contains the proof of Theorem 5, while the remaining subsections address Theorems 6 and 7.

6.1. An extension of the known interval for $K_{2,t}$.

Proof of Theorem 5. Let $K$ be a $p$-core CRG for $p \in \left[\frac{2}{t+1},1\right]$ that does not permit $K_{2,t}$ embedding for $t \geq 5$. If we assume that $g_K(p) < g_{K(t-1)}(p) = (1-p)/(t-1)$, then by Theorem 19 $K$ has only black vertices and no black edges.

Again, we partition the vertices of $K$ into three sets $\{u_0\}, U = \{u_1, \ldots, u_t\}$ and $W$, where $u_0$ is a fixed vertex with maximum weight $x$; $U$ is the set of all vertices in the gray neighborhood of $u_0$ with $u_1$ a vertex of maximum weight $x_1$ in $U$; and $W$ is the set of all remaining vertices, or those vertices adjacent to $u_0$ via white edges. Finally, let $d_G(u_i)$ signify the sum of the weights of all vertices in the gray neighborhood of $u_i$.

Then by Lemma 20, the total weight of the vertices in $W$ is at least

$$d_G(u_1) - (t-3)x_1 - x,$$

since no vertex in $U$ can be adjacent to more than $t-3$ other vertices in $U$ without forming a book $B_{t-2}$ gray subgraph with $u_0$. Thus,

$$x + d_G(u_0) + [d_G(u_1) - (t-3)x_1 - x] \leq 1.$$

Applying Proposition 23 and letting $g_K(p) = g$,

$$2 \left( \frac{p - g}{p} \right) + \frac{1-2p}{p} x + \left[ \frac{1-2p}{p} - (t-3) \right] x_1 \leq 1$$
$$2(p-g) - p + (1-2p)x \leq [p(t-1) - 1] x_1$$
$$2(p-g) - p + (1-2p)x \leq [p(t-1) - 1] x_1$$
$$p - 2g \leq [p(t+1) - 2] x.$$

Since $p \geq \frac{2}{t+1}$ and $x \leq \frac{g}{1-p}$ by Proposition 23

$$p - 2g \leq [p(t+1) - 2] \frac{g}{1-p}$$

$$\frac{1-p}{t-1} \leq g.$$

By Theorem 19 $g \leq \frac{1-p}{t-1}$, for $p \in \left[\frac{2}{t+1},1\right]$, so $ed_{\text{Forb}(K_{2,t})}(p) = \frac{1-p}{t-1}$. □

We will now show that this result is enough to determine the maximum value of $ed_{\text{Forb}(K_{2,t})}(p)$ for odd $t$. 

17
6.2. A construction for odd \( t \).

**Proposition 35.** Let \( \mathcal{H} = \text{Forb}(K_{2,t}) \) for odd \( t \). Then \( \text{ed}_\mathcal{H}(p) \leq 1/(t+1) \).

**Proof.** Let \( K \) be the CRG consisting of \( t+1 \) black vertices with white subgraph forming a perfect matching and all other edges gray. The CRG, \( K \), does not contain a gray \( K_{2,t} \) or book \( B_{t-2} \), and so by Lemma [20] \( K \) forbids a \( K_{2,t} \) embedding.

The CRG, \( K \), contains exactly \( (t+1)/2 \) white edges, so by Equation (3),

\[
\text{ed}_K(p) = \frac{1}{(t+1)^2} \left[ p \left( 2 \cdot \frac{t+1}{2} \right) + (1-p)(t+1) \right] = \frac{1}{t+1}.
\]

Therefore, \( \text{ed}_\mathcal{H}(p) \leq 1/(t+1) \). \( \square \)

Since by Theorem [3] \( \text{ed}_{\text{Forb}(K_{2,t})}(\frac{2}{t+1}) = \frac{1}{t+1} \), and by Proposition [35] \( \text{ed}_{\text{Forb}(K_{2,t})} \leq \frac{1}{t+1} \), we have that \( d^*_\text{Forb}(K_{2,t}) = \frac{1}{t+1} \) for odd \( t \geq 5 \).

6.3. A general lower bound for \( t \).

We conclude this section by determining a general lower bound for the edit distance function of \( \text{Forb}(K_{2,t}) \). It is the lower bound from Theorem [11] and it allows us to make the claim in Theorem [1] that, in the case of odd \( t \), there is a nondegenerate interval \( p^*_H \) that achieves the maximum value of the function.

**Proof of Theorem [11]** Here we use the standard bounds from Propositions [22] and [23]. Let \( g = g_K(p) \), where \( K \) is a black-vertex, \( p \)-core CRG, and let \( N_G(v) \) denote the gray neighborhood of a given vertex \( v \) in \( K \). Then if \( u_1, \ldots, u_\ell \) are the vertices in the gray neighborhood, \( U \), of a fixed vertex of maximum weight, \( u_0 \),

\[
\sum_{i=1}^\ell \left[ d_G(u_i) - x - x \left( N_G(u_i) \cap N_G(u_0) \right) \right] \leq (t-1)(1-x-d_G(u_0)).
\]

The left-hand side of this inequality calculates the weight of the total gray neighborhood of each vertex in \( U \) that must be contained in \( W \), the set of all vertices not in \( U \) or \( u_0 \). On the right-hand side we make use of the facts that \( x(W) = 1 - x - d_G(u_0) \) and that no vertex in \( W \) may be adjacent to more than \( t-1 \) vertices in \( U \) without violating Lemma [20] by forming a gray \( K_{2,t} \) with \( u_0 \). Thus, applying Proposition [22]

\[
\sum_{i=1}^\ell \left[ \frac{p-g}{p} - x + \frac{1-2p}{p} x(u_i) \right] - \sum_{i=1}^\ell x \left( N_G(u_i) \cap N_G(u_0) \right) \leq (t-1)(1-x-d_G(u_0)).
\]

Again considering Lemma [20] reveals that no vertex \( u_i \in U \) can have more than \( t-3 \) gray neighbors in \( U \) without inducing a gray book \( B_{t-2} \) with \( u_0 \). Therefore,

\[
\ell \left[ \frac{p-g}{p} - x \right] + \frac{1-2p}{p} d_G(u_0) - (t-3)d_G(u_0) \leq (t-1)(1-x-d_G(u_0))
\]

\[
\ell \left[ \frac{p-g}{p} - x \right] \leq (t-1)(1-x) - \frac{1}{p} d_G(u_0).
\]
Recalling that by Proposition 23 \( \frac{p - g}{p} \geq x \), we use the pigeon-hole bound from Fact 20 \( \ell \geq d_G(u_0)/x \) to get

\[
\frac{d_G(u_0)}{x} \left[ \frac{p - g}{p} - x \right] \leq (t - 1)(1 - x) - \frac{1}{p} d_G(u_0)
\]

\[
d_G(u_0) \left[ \frac{p - g}{p} - x \right] \leq (t - 1)x(1 - x) - \frac{x}{p} d_G(u_0)
\]

\[
d_G(u_0) \left[ \frac{p - g}{p} + \frac{1 - p}{p} x \right] \leq (t - 1)x(1 - x).
\]

By Proposition 22

\[
\left[ \frac{p - g}{p} + \frac{1 - 2p}{p} x \right] \left[ \frac{p - g}{p} + \frac{1 - p}{p} x \right] \leq (t - 1)x(1 - x).
\]

Collecting terms yields,

\[
\left( \frac{p - g}{p} \right)^2 + \left( \frac{p - g}{p} \right) \left( \frac{2 - 3p}{p} \right) - (t - 1) x + \left( \frac{1 - 2p}{p} \right) \left( \frac{1 - p}{p} \right) + (t - 1) \right] x^2 \leq 0,
\]

and so minimizing the left-hand side of the inequality with respect to \( x \), we have

\[
\left( \frac{p - g}{p} \right)^2 - \frac{(t - 1) - \left( \frac{p - g}{p} \right) \left( \frac{2 - 3p}{p} \right)}{4 \left( \frac{1 - 2p}{p} \right) \left( \frac{1 - p}{p} \right) + (t - 1)} \leq 0
\]

\[
\left( \frac{p - g}{p} \right)^2 (4t - 5) + 2 \left( \frac{p - g}{p} \right) (t - 1) \left( \frac{2 - 3p}{p} \right) - (t - 1)^2 \leq 0.
\]

Using the quadratic formula,

\[
p - \frac{g}{p} \leq \frac{2(t - 1) \left( \frac{2 - 3p}{p} \right) + \sqrt{4(t - 1)^2 \left( \frac{2 - 3p}{p} \right)^2 + 4(t - 1)^2(4t - 5)}}{2(4t - 5)}
\]

\[
p - \frac{g}{p} \leq \frac{t - 1}{4t - 5} \left[ 3p - 2 + \sqrt{(2 - 3p)^2 + (4t - 5)p^2} \right]
\]

\[
g \geq \frac{t - 1}{4t - 5} \left[ 3p - 2 + 2\sqrt{1 - 3p + (t + 1)p^2} \right].
\]

\[
\square
\]

The function in (11) achieves its maximum at \( p = \frac{2t - 1}{t(t + 1)} \), and that maximum is, in fact, \( \frac{1}{t + 1} \). Hence \( ed_{\text{Forb}(K_{2,t})}(p) \) is at least \( \frac{1}{t + 1} \) at \( p = \frac{2t - 1}{t(t + 1)} \) and is at least \( \frac{1}{t + 1} \) at \( p = \frac{2}{t + 1} \). As a result of concavity,

\[
ed_{\text{Forb}(K_{2,t})}(p) \geq \frac{1}{t + 1} \quad \text{for} \quad p \in \left[ \frac{2t - 1}{t(t + 1)}, \frac{2}{t + 1} \right].
\]

Equality holds whenever \( t \) is odd because, in that case, Proposition 35 gives that \( ed_{\text{Forb}(K_{2,t})}(p) \leq 1/(t + 1) \), so \( p_{\text{opt}}^t \) must be an interval. This concludes the proof of Theorem 7.

If \( t \geq 5 \), then we can analyze the first and second derivatives, with respect to \( p \), of

\[
p(1 - p) - \left( p - \frac{t - 1}{4t - 5} \left[ 3p - 2 + 2\sqrt{1 - 3p + (t + 1)p^2} \right] \right).
\]

The maximum difference between \( p(1 - p) \) and the lower bound in Theorem 6 on the interval \([0, \frac{2}{t + 1}]\) is \( \frac{1}{t + 1} \) and occurs when \( p = \frac{2t - 1}{t(t + 1)} \). We can also see that (6) is bounded below by \( \left( \frac{1}{t} - \frac{1}{t + 1} \right) p(1 - p) \).
7. Upper bound constructions

That we have been able to determine the entire edit distance function for Forb($K_{2,3}$) and Forb($K_{2,4}$) raises the question of whether it might be possible to do something similar for Forb($K_{2,t}$) when $t \geq 5$. That is, can we always find a few simple upper bounds that determine the entire edit distance function? In this section we show that when $t \geq 5$, the number and types of known upper bounds for the function increases significantly, though this does not necessarily preclude the possibility that a few, yet to be discovered, CRG constructions could determine the entire function.

In Section 7.1, we revisit the work in Section 5.1 on strongly regular graphs. In Section 7.2, we give some constructions inspired by the analysis of triangle-free graphs in [5].

7.1. Results from strongly regular graph constructions.

Recall that a strongly regular graph with parameters $(k, d, \lambda, \mu)$ is a $d$-regular graph on $k$ vertices such that each pair of adjacent vertices has $\lambda$ common neighbors, and each pair of nonadjacent vertices has $\mu$ common neighbors. Here we develop a function based on the existence of a strongly regular graph.

Suppose that $K$ is a CRG with all vertices black and all edges white or gray that is derived from a $(k, d, \lambda, \mu)$-strongly regular graph so that the edges of the strongly regular graph correspond to gray edges of $K$. In such a case we recall from Section 5.1 that

$$f_{S_k,d,\lambda,\mu}(p) = \frac{1}{k} + \left(\frac{k - d - 2}{k}\right)p.$$

As is commonly known (see [15], for instance), if a strongly regular graph with parameters $(k, d, \lambda, \mu)$ exists then it is necessary, though not sufficient, for

$$d(d - \lambda - 1) = \mu(k - d - 1).$$

If we substitute $\lambda = t - 3$ and $\mu = t - 1$ in this equation and then solve for $k$, we find that

$$k = \frac{t - 1 + d(d + 1)}{t - 1},$$

and substituting these values into $f_{S_k,d,\lambda,\mu}(p)$ yields

$$f_{S_k,d,\lambda,\mu}(p) = \frac{t - 1}{t - 1 + d(d + 1)} + \left(1 - \frac{(d + 2)(t - 1)}{t - 1 + d(d + 1)}\right)p.$$

Fixing $p$ and minimizing $f_{S_k,d,\lambda,\mu}(p)$ with respect to $d$ gives the following expression:

$$p\left(\frac{t - 2}{4t - 5}\right) + \frac{2(t - 1)}{4t - 5}\sqrt{1 - 3p + (t + 1)p^2},$$

which is equal to the lower bound from (1) in Theorem 6.

Of course, in order to even have a chance of actually attaining (7) with a strongly regular graph construction, both $d$ and $k = \frac{t - 1 + d(d + 1)}{t - 1}$ must be integers. This equation, however, provides something of a best case scenario for strongly regular graphs, and if there is a CRG, $K$, derived from a $(k, d, t - 3, t - 1)$-strongly regular graph that realizes equation (7), then $f_K(p)$ is tangent to the lower bound in (1) at

$$p = \frac{2d + 1}{(d + 1)(d + 3) - t},$$

determining the value of $ed_{\text{Forb}(K_{2,t})}(p)$ exactly.

The remaining upper bounds in Theorem 12 are the result of checking constructions from the known strongly regular graphs listed at [6]. Figure 4 (see Appendix A) is a chart of the relevant parameters and $f_K(p)$ functions for $5 \leq t \leq 8$. 
To our knowledge, it is not known whether, for fixed \( t \), there are a finite or infinite number of \((k,d,t-3,t-1)\)-strongly regular graphs. See Elzinga [9] for values of \( \lambda \) and \( \mu \) for which the number of strongly regular graphs with parameters \((k,d,\lambda,\mu)\) is known to be finite or infinite.

There is an additional construction defining the upper bound for \( t = 8 \) in Theorem [11] described in the following section.

7.2. Cycle construction.

**Definition 36** ([15], p.296). For two vertices \( x, y \in V(G) \), where \( G \) is a simple connected graph, let \( \text{len}(x,y) \) denote the length of the minimum path from \( x \) to \( y \). The \( r \)th power of \( G \), \( G^r \), is the graph with vertex set \( V(G^r) = V(G) \), and edge set \( E(G^r) = \{xy: x \neq y \text{ and len}(x,y) \leq r\} \).

Let \( C_k^t \) be the cycle on \( k \) vertices raised to the \( r \)th power. Define \( C_{k,r} \) to be the CRG on \( k \) black vertices with white edges corresponding to those in \( C_k^r \) and gray edges corresponding to those in the complement of \( C_k^r \). Recall that EW denotes the set of white edges for a given CRG. Then

\[
 f_{C_k^t}(p) = \frac{1}{k^2} [(1-p)k + 2p |\text{EW}|] = \frac{1}{k^2} [(1-p)k + 2p (rk)] = \left(\frac{2r-1}{k}\right)p + \frac{1}{k}.
\]

**Proposition 37.** \( C_{5+t,2} \) forbids a \( K_{2,t} \) embedding, and therefore \( ed_{\text{Forb}(K_{2,t})}(p) \leq \frac{3p+1}{5+t} \).

**Proof.** First, we check that \( C_{5+t,2} \) does not contain a gray \( K_{2,t} \). If \( u_1 \) and \( u_2 \) are any two vertices in \( C_{5+t}^2 \), then \(|(N(u_1) \cup N(u_2)) - \{u_1, u_2\}| \geq 4\). This inequality is justified by observing that two vertices \( u_1 \) and \( u_2 \) that are neighbors in \( C_{5+t}^2 \) have the smallest possible number of total neighbors in \( C_{5+t}^2 \), and this common neighborhood has order 4. It then follows that \(|N(u_1) \cap N(u_2)| \leq t - 1\) in the complement of \( C_{5+t}^2 \). Thus, \( C_{5+t,2} \) does not contain a gray \( K_{2,t} \).

Second, we check that \( C_{5+t,2} \) does not contain a gray \( B_{t-2} \). If \( u_1 \) and \( u_2 \) are any two nonadjacent vertices in \( C_{5+t}^2 \), then \(|(N(u_1) \cup N(u_2)) - \{u_1, u_2\}| \geq 6\). Therefore, by reasoning similar to above, \(|N(u_1) \cap N(u_2)| \leq t - 3\) in the complement of \( C_{5+t}^2 \), implying \( C_{5+t,2} \) does not contain a gray \( B_{t-2} \).

Thus, by Lemma [20] \( C_{5+t,2} \) forbids a \( K_{2,t} \) embedding, and therefore \( ed_{\text{Forb}(K_{2,t})}(p) \leq f_{C_{5+t,2}}(p) = \frac{3p+1}{5+t} \).

While there are several other orders and powers of cycles that would also lead to a construction forbidding \( K_{2,t} \) embedding, none of them have a corresponding \( f_K(p) \) value that beats the upper bound \( \min\{p(1-p), \frac{3p+1}{5+t}, \frac{1-p}{t+1}\} \), so we restrict our interest to this one.

For \( t \geq 5 \), \( f_{C_{5+t,2}}(p) \) is always an improvement on the bound \( \min\{p(1-p), \frac{1-p}{t+1}\} \) from Theorem [19] though it is improved upon or made irrelevant by bounds from strongly regular graphs, for \( t \leq 7 \). When \( t = 4 \), the function \( f_{C_{5,2}}(p) \) is tangent to the edit distance function at \( p = 1/3 \), where the edit distance function achieves its maximum value.

7.3. Füredi constructions.

As is observed in Lemma [20] and used in the exploration of the past two constructions, graphs that forbid \( K_{2,t} \) and \( B_{t-2} \) as subgraphs are of interest when looking for CRGs that forbid \( K_{2,t} \) embedding. The following results come from examining the bipartite versions of \( K_{2,t} \)-free graph constructions described by Füredi [10]. This strategy mimics the one used in [11] with Brown’s \( K_{3,3} \)-free construction.
Proof of Theorem 39. We take the construction described in [11] for a $K_{2,t}$-free graph $G$ on $n = (q^2 - 1)/(t - 1)$ vertices, each with degree $q$, where $q$ is a prime power so that $t - 1$ divides $q - 1$. We should note here that in the original construction from [11], loops were omitted, reducing the degree of some vertices to $q - 1$. It is to our advantage, however, to leave the loops in so that the final construction will be $q$-regular. By the same proof as in [11], the graph with loops still retains the property that no two vertices have a common neighborhood greater than $t - 1$ even when a looped vertex is considered to be in its own neighborhood.

Next, we create a CRG $K$, by taking two copies of the vertex set $\{v_1, \ldots, v_n\}$ from the $K_{2,t}$-free graph with loops described above: $\{v_1', \ldots, v_n'\}$, $\{v_1'', \ldots, v_n''\}$. Color all of these $k = 2n$ vertices black, and let $EG(K) = \{v_i'v_j'' : v_iv_j \in E(G)\}$ with all edges not in $EG(K)$ white.

The gray subgraph of $K$ is bipartite, so it cannot contain a $B_{t-2}$, and since no two vertices $v_i$ and $v_j$ from the original construction have more than $t - 1$ common neighbors, the common neighborhood of two vertices in the gray subgraph of $K$ is also at most $t - 1$. Thus by Lemma 20, $K$ forbids a $K_{2,t}$ embedding.

The CRG, $K$, has $k = 2n = (q^2 - 1)/(t - 1)$ vertices and $q(q^2 - 1)/(t - 1)$ gray edges, so by equation (3), $f_K(p)$ is as described in the statement of Theorem 8.

Remark 38. Though the property of being bipartite is sufficient to exclude a $B_{t-2}$ subgraph, using a bipartite $K_{2,t}$-free construction may not be the optimal choice. A more efficient CRG may be constructed from another graph that has a gray subgraph that is both $K_{2,t}$- and $B_{t-2}$-free, but, for instance, still contains triangles.

Nevertheless, we can discover more about the potential for these constructions to improve upon the bounds for $ed_{\text{Forb}(K_{2,t})}(p)$ by fixing $p$ and considering the general formula in Theorem 8 as a continuous function with respect to $q$.

Lemma 39. Let $t \geq 3$, and let $q_0 < q$ be prime powers such that $t - 1$ divides both $q_0 - 1$ and $q - 1$. If the CRG, $K_0$, is constructed according to the proof of Theorem 8 with parameter $q_0$ and if the CRG, $K$, is constructed according to the proof of Theorem 8 with parameter $q$, then $f_{K_0}(p) \leq f_K(p)$ for $p \in \left[\frac{2}{4+q_0}, \frac{1}{3}\right]$.

Proof. We begin the proof by fixing $p$ and $t$ and analyzing $\phi(q) = \frac{t - 1 + p(2q^2 - q(t - 1) - 2t)}{2(q^2 - 1)}$. Note that $f_{K_0}(p) = \phi(q_0)$, and $f_K(p) = \phi(q)$. Consider when the derivative

$$\phi'(q) = \frac{(t - 1)(q^2p + p + 4qp - 2q)}{2(q^2 - 1)^2}$$

is positive and, therefore, $\phi$ is increasing. Since the greater value of $q$ that makes $q^2p + p + 4qp - 2q = 0$ (note that the leading term is nonnegative) occurs at $q = \frac{(1 - 2p)(1 - 2p)^2 - p^2}{p}$, it follows that $\phi'(q) \geq 0$ when $q \geq \frac{(1 - 2p)(1 - 2p)^2 - p^2}{p}$. If $p < 1/3$ and $q_0 \geq \frac{2(1 - 2p)}{p}$, then

$$q > q_0 \geq \frac{2(1 - 2p)}{p} > \frac{(1 - 2p) + \sqrt{(1 - 2p)^2 - p^2}}{p}.$$

Thus, $\phi'(q) \geq 0$ for $\frac{2}{4+q_0} \leq p < 1/3$. Therefore $f_{K_0}(p) \leq f_K(p)$ for $p$ in this interval.

Additionally, we can make some statements about when we can expect constructions that originate from the $K_{2,t}$-free graphs described by Füredi [11] to improve upon the bound $p(1 - p)$ for any $q$.

Lemma 40. Fix $t \geq 9$, and let $q$ be a prime power such that $t - 1$ divides $q - 1$. Let $K$ be the CRG with parameter $q$ described in the proof of Theorem 8, hence $f_K(p) = \frac{t - 1 + p(2q^2 - q(t - 1) - 2t)}{2(q^2 - 1)}$. Then
for any sufficiently large prime power $q$ and corresponding $K$, there is an interval of values of $p$ on which $f_K(p) < p(1-p)$. Moreover as $q \to \infty$ the left-hand endpoints of these open intervals approach 0.

That is, we can find an infinite sequence of CRG constructions that improve upon the known bounds for $\text{Forb}(K_{2,t})$ when $t \geq 9$, and the intervals on which these improvements occur get arbitrarily close to 0.

**Proof.** We begin by observing that $f_K(p) = \frac{p(q(t-1) + 2t - 1) - (t-1)}{2(q^2-1)} = p - \frac{p(q(t-1) + 2t - 1) - (t-1)}{2(q^2-1)}$. Thus if $f_K(p) < p(1-p)$,

$$p - \frac{p(q(t-1) + 2t - 1) - (t-1)}{2(q^2-1)} < p - p^2$$

$$2p^2(q^2 - 1) < p(q(t-1) + 2(t-1)) - (t-1)$$

(8) $$2p^2(q^2 - 1) - p(t-1)(q + 2) + (t-1) < 0.$$

The minimum value of $2p^2(q^2 - 1) - p(t-1)(q + 2) + (t-1)$ occurs when $p = \frac{(t-1)(q+2)}{4(q^2-1)}$. Therefore, the inequality above is satisfied for some $q$ and $p$ values if and only if

$$2 \left[ \frac{(t-1)(q+2)}{4(q^2-1)} \right]^2 (q^2 - 1) - \left[ \frac{(t-1)(q+2)}{4(q^2-1)} \right] (t-1)(q + 2) + (t-1) < 0$$

$$(t-1) \left( 1 - \frac{(t-1)(q+2)^2}{8(q^2-1)} \right) < 0.$$

That is, $f_K(p)$ from the constructions in [10] is less than $p(1-p)$ for some value of $p$ if and only if $1 - \frac{(t-1)(q+2)^2}{8(q^2-1)} < 0$. For positive $q$, it is always the case that $(q + 2)^2 > q^2 - 1$, and so any $q$ satisfying the constraints of the original construction will improve upon the upper bound established by $p(1-p)$ for some $p$ when $t \geq 9$. Furthermore, for a fixed prime power $q$ for which $t-1$ divides $q-1$, it is a definite improvement for some open neighborhood around $p = \frac{(t-1)(q+2)}{4(q^2-1)}$. This value approaches 0 as $q \to \infty$, and there are an infinite number of prime powers $q$ such that $t-1$ divides $q-1$ (see [10]). Thus, it is the case that for arbitrarily small $p$, we can find some $q$ such that $f_K(p) < p(1-p)$.

□

Lemma [11] is an analysis of the Füredi constructions when $t \leq 8$. We then show that our bounds from these constructions do not have an effect on the value of $e_{\text{Forb}(K_{2,t})}(p)$ for $t \leq 8$.

**Lemma 41.** Fix $5 \leq t \leq 8$, and let $q$ be a prime power such that $t-1$ divides $q-1$. Let $K$ be the CRG with parameter $q$ described in the proof of Theorem [8] hence $f_K(p) = \frac{t-1+p(2q^2-q(t-1)-2t)}{2(q^2-1)}$. Then

$$q < \frac{(t-1) + \sqrt{(t-1)^2 + (9-t)(t+1)}}{\frac{1}{2}(9-t)}.$$

(9)

**Proof.** Returning to inequality [8] and performing a similar analysis to that in the proof of Lemma [11] we see that if $t \leq 8$, then $2p^2(q^2 - 1) - p(t-1)(q + 2) + (t-1) < 0$ for some value of $p$ if and only if

$$\frac{(t-1) - \sqrt{(t-1)^2 + (9-t)(t+1)}}{\frac{1}{2}(9-t)} < q < \frac{(t-1) + \sqrt{(t-1)^2 + (9-t)(t+1)}}{\frac{1}{2}(9-t)}.$$

The lower bound for $q$ described above is immaterial since for $t \leq 8$ it is always negative. The upper bound completes the proof of Lemma [11] □
Using Lemma 41 we generated the following table of possible $q$ values that obey the inequality in (9). Since we have already determined the entire edit distance function for $\text{Forb}(K_{2,3})$ and $\text{Forb}(K_{2,4})$, only $t = 5, 6, 7, 8$ needed to be considered:

| $t$ | possible $q$ values |
|-----|----------------------|
| 5   | 5                    |
| 6   | none                 |
| 7   | 7, 13                |
| 8   | 8, 29                |

A case analysis of the $f_K(p)$ functions corresponding to these $q$ values finds no improvement to the bounds established by $\min\{p(1-p), \frac{3p+1}{t+1}, \frac{1-p}{t-1}\}$, except in the cases when $t = 7$ and $q = 13$, and $t = 8$ and $q = 29$. In these cases, we see an improvement for the approximate ranges $p \in (0.125, 0.1358)$ and $p \in (0.0625, 0.06667)$, respectively, but even these improvements are surpassed by results from strongly regular graph constructions.

8. Conclusion

- Although we determine all of $ed_{\text{Forb}(K_{2,4})}(p)$, convexity allows $d^*_{\text{Forb}(K_{2,4})}$ to be determined with only Lemma 41. Furthermore, the generalized quadrangle $\text{GQ}(2,2)$ was unnecessary to compute this quantity, since $p(1-p) < 2/9$ for $p \in [0,1/3]$, and $p(1-p)$ is an upper bound for every function $ed_{\text{Forb}(K_{2,t})}(p)$, $t \geq 2$.
- Proposition 43 gives a nontrivial lower bound for a black-vertex, $p$-core CRG, $K$, that forbids a $K_{2,4}$ embedding. If we take inequality (5) and solve for $g$, then we see that, if $p \in (0,1/2)$, then
  \[ g_K(p) \geq \frac{2p + 6 - 6\sqrt{1 - 3p + 5p^2}}{11}, \]
  which is strictly larger than $p(1-p)$ for $p \in (0, (9 - 4\sqrt{3})/11)$ and corresponds to the general lower bound described in Theorem 43. In particular, there is a positive gap between the $g_K(p)$ functions for black-vertex CRGs and the CRG with one white and one black vertex.
- Ed Marchant reports having also proven that $ed_{\text{Forb}(K_{2,3})}(p) = \min\{p(1-p), (1-p)/2\}$, using different methods.
- While we have yet to determine the entire edit distance function $ed_{\text{Forb}(K_{2,4})}(p)$ when $t \geq 5$, strongly regular graph constructions have the potential to determine its value exactly, at least for certain values of $p$. It is likely that with the development of knowledge of strongly regular graphs, we will see improved upper bounds for $ed_{K_{2,4}}(p)$. Already, they provide significant improvements to previously known upper bounds, realizing the function value exactly in some cases.

- For $t \geq 9$, Füredi’s $K_{2,t}$-free construction leads to improvements to the bound $p(1-p)$ for values of $p$ arbitrarily close to 0. In fact, analyzing inequality (8) with respect to $q$, indicates that for fixed $p$ and $t \geq 9$, if an appropriate $q$ exists such that
  \[ \frac{t - 1 - \sqrt{(t-1)^2 - 8(t-1)(1-2p) - 16p^2}}{4p} < q < \frac{t - 1 + \sqrt{(t-1)^2 - 8(t-1)(1-2p) - 16p^2}}{4p}, \]
  then there is an improvement.

  Meanwhile, for $t \leq 8$, the upper bounds from these constructions, at least when the corresponding $f_K$ functions are considered, are inferior to those from alternative constructions. It is unknown, whether or not these improvements are best possible, or if there is another construction that could render them irrelevant too.
9. Acknowledgements

Special thanks to Sung-Yell Song for his knowledge and expertise on strongly regular graphs and Maria Axenovich for providing references on Füredi’s constructions. Thank you to Ed Marchant and Andrew Thomason for helpful conversations and to Stephen Hartke for providing some references on strongly regular graphs. We also wish to acknowledge some helpful comments from an anonymous referee.

References

[1] Noga Alon and Uri Stav, What is the furthest graph from a hereditary property?, Random Structures Algorithms 33 (2008), no. 1, 87–104.
[2] Maria Axenovich, André Kézdy, and Ryan Martin, On the editing distance of graphs, J. Graph Theory 58 (2008), no. 2, 123–138.
[3] József Balogh and Ryan Martin, Edit distance and its computation, Electron. J. Combin. 15 (2008), no. 1, Research Paper 20, 27.
[4] Christian Borgs, Jennifer Chayes, László Lovász, Vera T. Sós, Balázs Szegedy, and Katalin Vesztergombi, Graph limits and parameter testing, STOC’06: Proceedings of the 38th Annual ACM Symposium on Theory of Computing, ACM, New York, 2006, pp. 261–270.
[5] Stephan Brandt, On the structure of dense triangle-free graphs, Combin. Probab. Comput. 8 (1999), no. 3, 237–245.
[6] Andries E. Brouwer, Parameters of Strongly Regular Graphs. http://www.win.tue.nl/~aeb/graphs/srg/srgtab.html (accessed July 1, 2010).
[7] W. G. Brown, On graphs that do not contain a Thomsen graph, Canad. Math. Bull. 9 (1966), 281–285.
[8] Guantao Chen, Ralph J. Faudree, and Ronald J. Gould, Saturation numbers of books, Electron. J. Combin. 15 (2008), no. 1, Research Paper 118, 12.
[9] Randall J. Elzinga, Strongly regular graphs: values of λ and µ for which there are only finitely many feasible \((v, k, λ, µ)\), Electron. J. Linear Algebra 10 (2003), 232–239 (electronic).
[10] Zoltán Füredi, New asymptotics for bipartite Turán numbers, J. Combin. Theory Ser. A 75 (1996), no. 1, 141–144.
[11] Edward Marchant and Andrew Thomason, Extremal graphs and multigraphs with two weighted colours, Fete of combinatorics and computer science, Bolyai Soc. Math. Stud., vol. 20, János Bolyai Math. Soc., Budapest, 2010, pp. 239–286.
[12] Ryan Martin, The edit distance function and symmetrization. arXiv:1007.1897v4.
[13] Alexander Sidorenko, Boundedness of optimal matrices in extremal multigraph and digraph problems, Combinatorica 13 (1993), no. 1, 109–120.
[14] Andrew Thomason, Graphs, colours, weights and hereditary properties. Preprint.
[15] Douglas B. West, Introduction to Graph Theory, 2nd, Prentice Hall, Upper Saddle River, NJ, 2001.
Appendix A. Figures

| $t$ values | parameters | $f_K(p)$ |
|------------|------------|----------|
| $t \geq 5$ | (13, 6, 2, 3) | $(1 + 5p)/13$ |
|           | (40, 12, 2, 4) | $(1 + 26p)/40$ |
|           | (96, 19, 2, 4) | $(1 + 75p)/96$ |
| $t \geq 6$ | (10, 6, 3, 4) | $(1 + 2p)/10$ |
|           | (17, 8, 3, 4) | $(1 + 7p)/17$ |
|           | (26, 10, 3, 4) | $(1 + 14p)/26$ |
|           | (85, 20, 3, 5) | $(1 + 63p)/85$ |

| $t$ values | parameters | $f_K(p)$ |
|------------|------------|----------|
| $t \geq 7$ | (16, 9, 4, 6) | $(1 + 5p)/16$ |
|           | (36, 14, 4, 6) | $(1 + 20p)/36$ |
|           | (49, 16, 3, 6) | $(1 + 31p)/49$ |
|           | (64, 18, 2, 6) | $(1 + 44p)/64$ |
|           | (100, 22, 0, 6) | $(1 + 76p)/100$ |
|           | (156, 30, 4, 6) | $(1 + 124p)/156$ |
| $t \geq 8$ | (25, 12, 5, 6) | $(1 + 11p)/25$ |
|           | (76, 21, 2, 7) | $(1 + 53p)/76$ |
|           | (125, 28, 3, 7) | $(1 + 95p)/125$ |

Figure 4. Above are the known parameters (see [8]) and $f_K(p)$ functions from strongly regular graphs that provide an improvement upon the known upper bound for $ed_{\mathcal{H}}(p)$ for some interval of $p$ values, where $\mathcal{H} = \text{Forb}(K_{2,t})$ for $5 \leq t \leq 8$. Parameters with resulting bounds surpassed by other strongly regular graph constructions are omitted.

Figure 5. Upper and lower bounds (in solid and dashed respectively) for $ed_{\text{Forb}(K_{2,5})}(p)$. Points indicate tangency.

Figure 6. Difference between upper and lower bounds for $ed_{\text{Forb}(K_{2,5})}(p)$. 
Figure 7. Upper and lower bounds (in solid and dashed respectively) for $ed_{Forb(K_{2,6})}(p)$. Points indicate tangency.

Figure 8. Difference between upper and lower bounds for $ed_{Forb(K_{2,6})}(p)$.

Figure 9. Upper and lower bounds (in solid and dashed respectively) for $ed_{Forb(K_{2,7})}(p)$. Points indicate tangency.

Figure 10. Difference between upper and lower bounds for $ed_{Forb(K_{2,7})}(p)$.

Figure 11. Upper and lower bounds (in solid and dashed respectively) for $ed_{Forb(K_{2,8})}(p)$. In this instance there are no points of tangency.

Figure 12. Difference between upper and lower bounds for $ed_{Forb(K_{2,8})}(p)$. 

27
Appendix B. Proofs of Lemma 31 and Proposition 32

Proof of Lemma 31:

Proof. We break this into two cases: when $K$ does and does not have a gray triangle.

Case 1: $K$ has a gray triangle.

Let the gray subgraph of $K$ contain a triangle whose vertices are $v_1$, $v_2$ and $v_3$ with optimal weights $y_1$, $y_2$ and $y_3$, respectively. Because $K$ has no gray $B_2$, we know that no pair of the vertices $v_1, v_2, v_3$ have a common gray neighbor other than the remaining vertex in the triangle. Letting $g = g_K(p)$, we have the following because the sum of the optimal weights on all vertices in $K$ is 1:

$$y_1 + y_2 + y_3 + \frac{\sum_{i=1}^{3} [d_G(v_i) - (y_1 + y_2 + y_3 - y_i)]}{3} \leq 1.$$

Then, applying Proposition 32

$$y_1 + y_2 + y_3 + 3 \left( \frac{p - g}{p} \right) + \frac{1 - 2p}{p} (y_1 + y_2 + y_3 - 2(y_1 + y_2 + y_3) \leq 1$$

and so

$$\frac{2p - 3g}{p} \leq \left( \frac{3p - 1}{p} \right) (y_1 + y_2 + y_3) \leq \frac{3p - 1}{p}.$$

Consequently, $g \geq (1 - p)/3$ with equality if and only if $y_1 + y_2 + y_3 = 1$; i.e., $K$ itself is a gray triangle.

Case 2: $K$ has no gray triangle.

Let $u_0$ be a vertex of largest weight, $x = x(u_0)$, and let $U = N_G(u_0)$. The absence of a gray triangle means that there are no gray edges between pairs of vertices in $U$. Furthermore, no vertex in $W$ can be adjacent to more than three vertices in $U$ via a gray edge, since by Lemma 20, the gray subgraph of $K$ does not contain a $K_{2,4}$.

Let $u_1, \ldots, u_\ell$ be an enumeration of the vertices in $U$ with weights $x_1, \ldots, x_\ell$, respectively, and $g = g_K(p)$. Then

$$\sum_{i=1}^{\ell} (d_G(u_i) - x) \leq 3x(W)$$

$$\leq 3(1 - x - x(U)),$$

and applying Proposition 32 to compute $d_G(u_i)$,

$$\sum_{i=1}^{\ell} \left( \frac{p - g}{p} + \frac{1 - 2p}{p} x_i - x \right) \leq 3(1 - x - x(U))$$

$$\ell \left( \frac{p - g}{p} - x \right) + \frac{1 - 2p}{p} x(U) \leq 3(1 - x) - 3x(U)$$

(10) $$\ell \left( \frac{p - g}{p} - x \right) \leq 3(1 - x) - \frac{1 + p}{p} x(U).$$
First, suppose $\ell \geq 5$. Then, from inequality (10), we have
\[
5 \left( \frac{p - g}{p} - x \right) \leq 3(1 - x) - \frac{1 + p}{p} x(U),
\]
and applying Proposition 22 again,
\[
5 \left( \frac{p - g}{p} - x \right) \leq 3(1 - x) - \frac{1 + p}{p} \left( \frac{p - g}{p} + \frac{1 - 2p}{p} x \right)
\]
\[
\frac{1 + 6p}{p} \cdot \frac{p - g}{p} - 3 \leq \left( 5 - 3 - \frac{1 + p}{p} \cdot \frac{1 - 2p}{p} \right) x
\]
\[
p(1 + 3p) - g(1 + 6p) \leq x \left( 4p^2 + p - 1 \right).
\]
If $4p^2 + p - 1 < 0$, then we may use the fact that $x > 0$,
\[
g \geq \frac{p(1 + 3p)}{1 + 6p} = \frac{1 - p}{3} + \frac{(3p - 1)(1 + 5p)}{3(1 + 6p)}.
\]
If $4p^2 + p - 1 \geq 0$, then we use Proposition 23 and substitute $x = g/(1 - p)$,
\[
p(1 + 3p) - g(1 + 6p) \leq \frac{g}{1 - p} \left( 4p^2 + p - 1 \right)
\]
\[
p(1 + 3p) \leq g \left( \frac{6p - 2p^2}{1 - p} \right)
\]
\[
\frac{1 - p}{3} + \frac{(1 - p)(11p - 3)}{6(3 - p)} \leq g.
\]
Regardless of the value of $p \in (1/3, 1/2)$, if $\ell \geq 5$, then $g > (1 - p)/3$. Therefore, we may assume that $\ell \leq 4$.

Second, suppose $\ell \leq 2$. Then by Fact 26 we have $\ell \geq x(U)/x$, yielding
\[
x(U)/x \leq \ell \leq 2,
\]
and so bounding $x(U)$ using Proposition 22,
\[
\frac{1}{x} \left( \frac{p - g}{p} + \frac{1 - 2p}{p} x \right) \leq 2
\]
\[
\frac{p - g}{p} \leq \frac{4p - 1}{p} x.
\]
Using Proposition 23 $x \leq \frac{g}{1 - p}$ yields
\[
\frac{p - g}{p} \leq \frac{4p - 1}{p} \cdot \frac{g}{1 - p}
\]
\[
p(1 - p) \leq 3pg,
\]
and so if $\ell \leq 2$, then $g \geq (1 - p)/3$, with equality if and only if $x = g/(1 - p)$, and consequently, $K$ is a gray triangle. So, we may further assume that $\ell \in \{3, 4\}$.

Third, suppose $\ell = 3$. Then
\[
x(U)/x \leq 3
\]
\[
\frac{p - g}{p} \leq \frac{5p - 1}{p} x
\]
\[
\frac{p - g}{5p - 1} \leq x.
\]
Returning to inequality (10), we have
\[
3 \left( \frac{p - g}{p} - x \right) \leq 3(1 - x) - \frac{1 + p}{p} x(U)
\]
\[
\frac{1 + 4p}{p} \cdot \frac{p - g}{p} - 3 \leq - \left[ \frac{1 + p}{p} \cdot \frac{1 - 2p}{p} \right] x
\]
\[
p(1 + p) - g(1 + 4p) \leq -(1 + p)(1 - 2p) \left( \frac{p - g}{5p - 1} \right)
\]
\[
p(1 + p)(5p - 1) + p(1 + p)(1 - 2p) \leq g \left[ (1 + p)(1 - 2p) + (1 + 4p)(5p - 1) \right]
\]
\[
\frac{1 + p}{6} \leq g
\]
\[
\frac{1 - p}{3} + \frac{3p - 1}{6} \leq g.
\]
If \( \ell = 3 \), then \( g > \frac{1}{3} \).

Fourth, and finally, suppose \( \ell = 4 \). Then
\[
\frac{x(U)}{x} \leq 4
\]
\[
\frac{p - g}{p} \leq 6p - 1
\]
\[
\frac{p - g}{6p - 1} \leq x.
\]

Returning to inequality (10), we have
\[
4 \left( \frac{p - g}{p} - x \right) \leq 3(1 - x) - \frac{1 + p}{p} x(U)
\]
\[
\frac{1 + 5p}{p} \cdot \frac{p - g}{p} - 3 \leq \left[ 4 - 3 - \frac{1 + p}{p} . \frac{1 - 2p}{p} \right] x
\]
\[
p(1 + 2p) - g(1 + 5p) \leq \left[ 3p^2 + p - 1 \right] x.
\]
If \( 3p^2 + p - 1 < 0 \), then we use the fact that \( x \geq (p - g)/(6p - 1) \):
\[
p(1 + 2p) - g(1 + 5p) \leq \left[ 3p^2 + p - 1 \right] \left[ \frac{p - g}{6p - 1} \right]
\]
\[
p(1 + 2p) - \frac{p(3p^2 + p - 1)}{6p - 1} \leq g \left[ 1 + 5p - \frac{3p^2 + p - 1}{6p - 1} \right]
\]
\[
\frac{1 + 3p}{9} \leq g
\]
\[
\frac{1 - p}{3} + \frac{2(3p - 1)}{9} \leq g.
\]

If \( 3p^2 + p - 1 \geq 0 \), then we use Fact 25 to bound \( x \leq \frac{g}{1 - p} \),
\[
p(1 + 2p) - g(1 + 5p) \leq \left[ 3p^2 + p - 1 \right] \left[ \frac{g}{1 - p} \right]
\]
\[
p(1 + 2p) \leq g \left[ 1 + 5p + \frac{3p^2 + p - 1}{1 - p} \right]
\]
\[
\frac{(1 - p)(1 + 2p)}{5 - 2p} \leq g
\]
\[
\frac{1 - p}{3} + \frac{2(4p - 1)(1 - p)}{3(5 - 2p)} \leq g.
\]
Regardless of the value of \( p \in (1/3, 1/2) \), if \( \ell = 4 \), then \( g > (1 - p)/3 \).
This ends Case 2 and the proof of the lemma.

**Proof of Proposition 33:**

*Proof.* Let \( u_1, \ldots, u_\ell \) be an enumeration of the vertices of \( U \). Then

\[
\sum_{i=1}^{\ell} (d_G(u_i) - x) - x(U_1) \leq 3(1 - x - x(U)),
\]

and applying Proposition 22

\[
\sum_{i=1}^{\ell} \left( \frac{p - g}{p} + \frac{1 - 2p}{p} x(u_i) - x \right) - x(U_1) \leq 3(1 - x - x(U)).
\]

Simplification yields the first inequality. The second inequality results from observing that \( x(U_1) \leq x(U) \).