Statistical Inference of a canonical dictionary of protein substructural fragments

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Abstract

Proteins are biomolecules of life. They fold into a great variety of three-dimensional (3D) shapes. Underlying these folding patterns are many recurrent structural fragments or building blocks (analogous to ‘LEGO® bricks’). This paper reports an innovative statistical inference approach to discover a comprehensive dictionary of protein structural building blocks from a large corpus of experimentally determined protein structures. Our approach is built on the Bayesian and information-theoretic criterion of minimum message length. To the best of our knowledge, this work is the first systematic and rigorous treatment of a very important data mining problem that arises in the cross-disciplinary area of structural bioinformatics. The quality of the dictionary we find is demonstrated by its explanatory power – any protein within the corpus of known 3D structures can be dissected into successive regions assigned to fragments from this dictionary. This induces a novel one-dimensional representation of three-dimensional protein folding patterns, suitable for application of the rich repertoire of character-string processing algorithms, for rapid identification of folding patterns of newly-determined structures. This paper presents the details of the methodology used to infer the dictionary of building blocks, and is supported by illustrative examples to demonstrate its effectiveness and utility.

1 Introduction

Proteins are molecules central to life. They are responsible for biological and cellular functions in all organisms. Each protein folds into a three-dimensional (3D) shape, determined by the intrinsic properties of its sequence or chain of amino acid residues. Among the major triumphs of modern science are the techniques to experimentally determine the 3D structures of proteins at atomic resolution. Worldwide structure determination efforts have resulted in a fast growing public database, the protein data bank (wwPDB) [1]. Currently wwPDB contains atomic coordinates of 91,550 experimentally solved protein structures, whose size is doubling every five years. This database provides a rich source of structural and architectural information for knowledge discovery and data mining applications that contribute to the advances made in life sciences in medicine.

Understanding the architectural principles of protein 3D structure is fundamental to biological research. It is well known that protein folding patterns contain recurrent structural themes, commonly helices and pleated sheets [2 3]. However the identification of a canonical set of building blocks (analogous to LEGO® bricks) of protein structures still remains an important open questions in biology.
Previous investigations have sought to identify a dictionary of fragments as building blocks of protein structures [4, 5, 6, 7, 8, 9, 10]. (By fragment we mean a contiguous region within the folding pattern of a protein – this can be viewed as a 3D analogue of a 1D substring.) However, these approaches largely rely on ad hoc clustering of short fragments of some fixed-length, usually 4 to 10 amino acid residues long. The restriction of generating fixed-length fragment libraries is an artificial constraint, mainly employed to work around the difficulty of the search problem that manifests when trying to identify recurring fragments of arbitrary length. Thus, the question, what is the canonical dictionary of fragments (of arbitrary lengths) of which all proteins are made, remains fundamentally unsolved.

This paper addresses the above question by framing it as as a statistical inference problem. Our approach relies on the Bayesian method of minimum message length inference [11, 12], where the optimal fragment dictionary is defined objectively as the one which permits the most concise explanation (or technically, shortest lossless encoding) of the coordinates of a collection of source protein structures. To the best of our knowledge, our work is the first objective and systematic treatment of this important question, addressed using a statistically rigorous approach which investigates the compressibility of protein coordinate data.

We mine these building blocks from a collection of 8992 experimentally determined protein structures, whose coordinates were solved at atomic resolution, and available from the wwPDB [1]. These source structures are dissimilar in amino acid sequence to avoid experimental and selection bias observed within the wwPDB. In other words, the collection we use is comprehensive and unbiased, representative of all known protein folding patterns in the wwPDB.

Our approach discovered 1711 fragments or building blocks, ranging in length from 4 to 31 amino acid residues. This dictionary allows the efficient, lossless representation (or encoding) of any given protein structure. For a particular protein structure, the optimal lossless encoding contains a dissection (or segmentation) – that is, a designation of successive non-overlapping regions in the protein structures that match the assigned dictionary fragments – and a statement of spatial deviations (or corrections) that should be applied to the coordinates of each assigned dictionary fragment so that the coordinates corresponding to each region in the actual structure can be recovered losslessly. We note that the regions that do not efficiently match any dictionary fragment are assigned to a ‘null model’; in these cases, the spatial deviations bear the entire weight of the description – this is tantamount to stating the coordinates of the region in the source structure raw (or as is).

The organization of the paper is as follows. Section 2 gives a brief introduction to the minimum message length criterion. Section 3 provides the foundations of the dictionary inference problem using the MML framework. This involves designing transparent communication processes, developing lossless encoding strategies, and evolving efficient search algorithms. Sections 4 and 5 give these details. Sections 6 and 7 provide further technical details to compute the total explanation message length, and the strategy that allows us to identify the best dictionary of fragments. The paper concludes with Section 8 providing various results including illustrative examples of the effectiveness and the utility of the dictionary we discover.

## 2 Minimum Message Length Framework

Minimum Message Length (MML) [11, 12] is a hypothesis (or model) selection paradigm which links statistical inference with information theory and data compression.

MML is a Bayesian method of inference. Formally, let $E$ denote a mass of observed data (or evidence) and $H$ a hypothesis on the data. From Bayes’s theorem [13] we have: $P(H\&E) = P(H) \times P(E|H) = P(E) \times P(H|E)$, where $P(H)$ is the prior probability of hypothesis $H$, $P(E)$ is the prior probability of data $E$, $P(H|E)$ is the posterior probability of $H$ given $E$, and $P(E|H)$ is the likelihood.

In the Bayesian framework, two competing hypotheses can be compared using the ratios of their
posterior probabilities:

\[
\frac{P(H_1|E)}{P(H_2|E)} = \frac{P(H_1)P(E|H_1)}{P(H_2)P(E|H_2)}
\]

Usually, the goal of inference is to select the hypothesis with the highest posterior probability.

MML offers a complementary view of Bayesian inference by linking the probability of an event with the message length required to transmit (or communicate, explain, describe) it losslessly. The mathematical theory of communication [14] gives the relationship between the message length \( I(e) \) to communicate an event \( e \) losslessly, and its probability \( P(e) \): \( I(e) = -\log P(e) \).\(^4\)

Therefore, by applying Shannon’s insight to Bayese theorem above, we get:

\[
I(H\&E) = I(H) + I(E|H) = I(E) + I(H|E)
\]

Similarly, two competing hypotheses can be compared as:

\[
I(H_1|E) - I(H_2|E) = I(H_1) + I(E|H_1) - I(H_2) - I(E|H_2)
\]

It directly follows that the best hypothesis \( H^* \) is the one for which the expression \( I(H^*) + I(E|H^*) \) takes the \textit{minimum} value. (Notice, this is equivalent to maximizing the posterior probability of the hypothesis given the data, \( P(H^*|E) \).)

MML is best understood as a communication process between an imaginary pair of transmitter (Alice) and receiver (Bob) connected by a Shannon channel. Alice’s objective is to send the data \( E \) using an explanation message in a form that Bob can receive it losslessly. Alice and Bob agree on a \textit{codebook} containing general rules of communication composed solely of common knowledge about typical, hypothetical data. Anything that is not a part of the codebook must be strictly transmitted as a part of the explanation message. If Alice can find the best hypothesis, \( H^* \) on the data, Bob will receive a decodable explanation message most economically.

Alice sends the explanation message of \( E \) in two parts. In the first part she transmits the best hypothesis, \( H^* \), she could find on the data \( E \) taking \( I(H^*) \) bits to communicate. In the second, Alice transmits the details of the observed data \( E \) not explained by the hypothesis \( H^* \), taking \( I(E|H^*) \) bits to communicate. (That is, this part correspond to the deviations of the observed data \( E \) with respect to \( H^* \).) Notice that MML inference gives a natural trade-off between hypothesis complexity (\( I(H^*) \)) and quality of its fit to the data (\( I(E|H^*) \)).

3 Inferring the dictionary using the MML criterion

\textbf{Preliminaries:} Let \( \mathcal{C} \) denote a \textit{collection} of source protein structures \( \{ \mathcal{P}_1, \mathcal{P}_2, \ldots, \mathcal{P}_{|\mathcal{C}|} \} \). In this work we use a subset of 8992 structures from the protein data bank after removing amino acid sequence bias. That is, no two structures in the collection \( \mathcal{C} \) have a sequence similarity greater than 40%.

Any protein structure \( \mathcal{P} \) is represented as an ordered list of \((x, y, z)\) coordinates of its alpha Carbon (\( C_\alpha \)) atoms along the protein backbone, denoted here as \( \mathcal{P} = \{ p_1, \ldots, p_{|\mathcal{P}|} \} \). All protein coordinates are defined in Angstrom units (\( 1\text{Å} = 10^{-10} \text{ meters} \)).

Let \( \mathcal{D} = \{ \mathcal{Q}_1, \mathcal{Q}_2, \ldots, \mathcal{Q}_{|\mathcal{D}|} \} \) denote a \textit{dictionary} of fragments. Each dictionary \textit{element} \( \mathcal{Q} = \{ q_1, \ldots, q_{|\mathcal{Q}|} \} \in \mathcal{D} \) is a substructural fragment (i.e., a list of coordinates corresponding to a consecutive region) derived from some \( \mathcal{P} \in \mathcal{C} \), of arbitrary length (\( |\mathcal{Q}| < |\mathcal{P}| \)).

We note that each \((x, y, z)\) comes specified (in the protein data bank) to 3 positions after the decimal place. Since we are dealing with inference based on lossless compression, we denote \( \epsilon = 0.001 \) as a parameter that specifies the accuracy to which coordinate data should be stated.

\(^4\)The unit of measurement of information depends on the base of the logarithm; \( \log_2 \) gives message lengths measured in bits, while \( \ln \) gives the same measured in nits.
Rationalizing this problem in the MML framework: In this work, any dictionary $D$ of fragments is a hypothesis of building blocks on a collection of structures $C$, with its observed coordinates acting as evidence for inference. Therefore, using the information-theoretic restatement of Bayes’s theorem describe in Section 2 we get:

$$I(D \& C) = I(D) + I(C|D).$$

Rationalizing Equation 3.1 as a communication process between Alice and Bob, the measure of quality of any proposed dictionary of substructures is the total length of the explanation message that Alice transmits to Bob so that all the coordinates in the collection of source structures are received losslessly. Given that even unrelated proteins contain common, recurrent fragments (or building blocks), Alice and Bob could reasonably hypothesize that they could apply this observation to transmit the coordinates of various structures more concisely, by using the dictionary of building blocks as the basis of communication. It is intuitive to see that the better a dictionary of fragments in terms of how well they describe (i.e., fit) the observed coordinates, the more economical is the description of the source structures in the collection.

To be useful for illuminating common building blocks over all proteins, the dictionary must be the same for all structures; that is, the dictionary does not change regardless of an individual structure that is being transmitted. Before transmitting the coordinates of the collection, Alice first encodes and sends Bob the canonical dictionary of fragments (taking $I(D)$ bits). With this information, Bob has a dictionary of substructures but not the coordinates of the source structures in the collection. Note again that Alice needs to send the dictionary only once (as a header (or first part) of the total explanation message): she need not restate it as part of the subsequent encodings of coordinates of particular protein structures being transmitted. Each subsequent message consists of the segmentation, that is, the optimal assignment of successive regions in a protein structure to dictionary fragments, plus the corrections (or vector deviations) required because each region in the source protein deviates from its assigned dictionary element. This takes $I(C|D)$ bits.

In proposing a dictionary Alice and Bob face a tradeoff. They could use a large, all-encompassing fragment dictionary whose elements fit regions of proteins very well, leaving only small deviations in the assigned regions of the proteins to be described. In this case the explanation message length for each protein would be dominated by the explanation cost of the dictionary and the assignment of dictionary fragments to regions, because a large dictionary requires a larger message to explain itself and to nominate a dictionary element. Alternatively, they could choose a small dictionary, in which case the message length would be dominated by the transmission of the corrections (vector deviations). As described in Section 2, the MML criterion provides an objective tradeoff between the dictionary complexity and its fit with the coordinate data observed in the collection.

*Optimality criterion:* The optimal dictionary involves finding a dictionary of fragments that minimizes the total message length equation shown in Equation 3.1. To achieve this involves the following criteria:

1. Assume some dictionary $D$ is given (containing arbitrary number of fragments, each of arbitrary length). According to the MML framework, the optimal encoding of a particular protein structure $P$ using the specified dictionary $D$ is defined as the combination, of minimal encoding length, of assignments of successive non-overlapping regions in $P$ to fragments in $D$, plus statements of spatial deviations relative to each assigned fragment per region to recover the observed coordinates in $P$ losslessly. Sections 4 and 5 present all the technical details to achieve this.

2. Next, given the above method to optimally encode a particular protein using a specified dictionary, the optimal encoding of a collection of protein structures $C$, all using the same fixed dictionary
\( \mathcal{D} \), requires the one-off statement of the dictionary \( \mathcal{D} \) (as a header to the subsequent explanation message), plus the optimal encodings of each individual protein \( P_i \in \mathcal{C} \) using the method in Step 1. Section 6 presents the details to achieve the optimal encoding of a collection with any specified dictionary.

3. Finally, given the method to optimal encode a collection of proteins in Step 2, an optimal dictionary for the collection of protein structures can be objectively defined as the one for which the one-off specification cost of the dictionary \( \mathcal{D} \), plus the sum of the optimal encodings of all the proteins \( P_i \in \mathcal{C} \), yields the shortest explanation message. Section 7 describes the search to infer the optimal dictionary.

4 Null model message length

The null model implies that the coordinates are being transmitted raw. However this message must not be wilfully inefficient. Recently, Konagurthu et al. [15] introduced statistical models to encode protein coordinate data, while dealing with a completely different problem of assigning secondary structure to protein coordinates. This work builds on their encoding schemes.

The null model encoding relies on the empirical observation that the distance between successive \( C_\alpha \) atoms in any protein chain is highly constrained at around 3.8 Å with a small deviation about this mean. For a chain (or ordered list) of coordinates \( \{p_1, p_2, \ldots, p_n\} \), any \( p_i \) can be transmitted raw given its previous coordinate \( p_{i-1} \) as follows: First the distance \( r_i \) between \( p_{i-1} \) and \( p_i \) is encoded and transmitted over a normal distribution \( \mathcal{N}(r_i; \mu, \sigma) \) stated to \( \epsilon \) accuracy, with \( \mu = 3.8 \) Å and \( \sigma = \pm 0.2 \) Å. This takes

\[
I_N(r_i) = \log_2 \left( \frac{\sqrt{2\pi}\sigma}{\epsilon} \right) + \frac{(r_i - \mu)^2}{2\sigma^2} \log_2 e \text{ bits.}
\]

With this information, the receiver Bob knows that \( p_i \) lies on a sphere of radius \( r_i \) centred at \( p_{i-1} \), but does not yet know its precise direction to pinpoint its location on the sphere. Assuming that \( p_i \)'s direction is potentially uniformly distributed (consistent with a null hypothesis), Alice can discretise the sphere’s surface into cells (using a common convention shared between Alice and Bob), each with area of \( \epsilon^2 \). This discretisation can then be used to transmit \( p_i \) as a cell number \( c_i \) over a variable length integer code. With the knowledge of \( p_{i-1}, r_i \) and \( c_i \), Bob can reconstruct \( p_i \) to the stated coordinate accuracy of \( \epsilon \) on each component. Thus, stating any cell number takes

\[
I_{\text{uniform}}(c_i) = -\log_2 \left( \frac{\epsilon^2}{4\pi r_i^2} \right) = \log_2 (4\pi r_i^2) - 2 \log_2 \epsilon \text{ bits.}
\]

Combining Equations 4.2 and 4.3, each coordinate \( p_i \) with respect to its previous coordinate \( p_{i-1} \) can be stated raw using this sphere approach in

\[
I_{\text{sphere}}(p_i) = I_N(r_i) + I_{\text{uniform}}(c_i) \text{ bits.}
\]

Building on the above, a chain of coordinates of any protein \( \mathcal{P}\{p_1, p_2, \ldots, p_{|\mathcal{P}|}\} \) can be transmitted as a null model message incrementally using this sphere approach.\footnote{Since the original frame of reference of protein coordinates have no special meaning, the coordinates can be translated such that the first coordinate \( p_1 \) lies on the origin and, hence, does not need to be transmitted explicitly as part of the message. Note, however, that even if \( p_1 \) is not assumed to be the origin, its encoding will add a constant cost to the total message length.} Alice sends Bob the length of the protein \( (|\mathcal{P}|) \) over a variable length integer code, followed by transmitting incrementally the coordinates
Transmitting $p_2$ given $p_1$, then $p_3$ given $p_2$ and so on. We note that the code length of transmitting any positive integer $n$ over a variable length $\log^* n$ distribution is:

\[
I_{\log^*}(n) = \log^*_2(n) + \log_2(2.865) \text{ bits}
\]

where $\log^*_2(n) = \log_2 n + \log_2 \log_2 n + \cdots$ (over all +ve terms).

Therefore, the null model message length to describe any protein chain takes

\[
I_{\text{null}}(p_1, \ldots, p_n) = I_{\log^*}(|P|) + \sum_{i=1}^{|P|} I_{\text{sphere}}(p_i) \text{ bits.}
\]

5 Optimal encoding of a particular protein structure given a dictionary for fragments

Let us assume that there exists a dictionary of fragments $\mathcal{D}$ which both Alice and Bob agree to use to communicate any chain of protein coordinates $P = \{p_1, \ldots, p_n\}$.

To encode the coordinates (i.e., compress the coordinates of $P$ using $\mathcal{D}$), Alice’s will dissect the protein chain into non-overlapping regions, each of which is assigned to some dictionary fragment. These assigned dictionary fragments are used as the basis to communicate the observed coordinates in the respective regions of the protein economically. This is possible if only if the assigned dictionary fragment has a good spatial fit with the region.

Implicit to the model of communication, in addition to the dictionary, is the null model. For any regions in the protein, the null model competes with various possible explanations using the dictionary fragments – if there is no compression to be gained (with respect to the null model message) using any one of the dictionary fragment to communicate some region, then the regions must be transmitted raw.

Thus, this leads to an interesting optimization problem: find the best dissection of proteins into (variable-length, non-overlapping) regions, each assigned to one of the fragments in the given dictionary or the null model. Before, we solve this problem optimally (see Section 5.4), we establish the procedure to encode the coordinates of a region using one of: (a) dictionary fragment (Section 5.1) or (b) null model (Section 5.2).

5.1 Describing a region using a dictionary fragment

The methodology to communicate of any region $\{p_i, \ldots, p_j\} \in P, 1 \leq i < j \leq (|P| \equiv n)$ using an assigned dictionary fragment $Q = \{q_1, \ldots, q_{|Q|}\} \in \mathcal{D}$ is described here. (Note, that a fragment can only be assigned provided the length of $Q$ is same as the length of the observed region, that is, $|Q| = j - i + 1$ residues.)

Broadly, the coordinates of the region are transmitted in two steps. First, the end point of the region $p_j$ is transmitted, followed by each of its interior points, $p_{i+1}, p_{i+2}, \cdots, p_{j-1}$.

Transmitting $p_j$. The end point $p_j$ is transmitted using the sphere approach similar to the one described in Section 4. The distance $d_{ij}$ between the start ($p_i$) and end ($p_j$) points is first transmitted using a normal distribution where the mean ($\mu$) is taken as the distance $d^*$ between the start ($q_1$) and end points ($q_{|Q|}$) of the assigned fragment. The standard deviation $\sigma$ of the end point is set to $\min((j-i) \times 0.2\AA, 3\AA)$ based on the length of the fragment, $|Q|$. Thus, using Equation 4.4 stating the radius and direction using the sphere approach takes $I_{\text{sphere}}(p_j)$ bits.

Transmitting interior points $p_{i+1}, \cdots, p_{j-1}$. The regions start and end points are now known to Bob, and there are $j-i-1$ interior points yet to be transmitted. Alice uses the following procedure to transmit the interior points using an assigned dictionary fragment. The coordinates in $Q$ are orthogonally transformed
interior points of the region, \( p \) and end \((p)\) circumference of a circle of radius \( r \) distance of an interior point within the dictionary fragment from its axis of rotation. Dividing the region, \( p \) observed region, \( |Q| \) the sum of the squared error of the interior points of the region with the corresponding interior points of the dictionary model is minimized.

This spatial transformation is related to the more general superposition problem that minimizes the sum of the squared distance between two corresponding vector sets which has an analytical solution [16]. However, the transformation here is constrained such that the first points of the two sets are the same (constraint \([a]\)) and the rotational axis for the dictionary model is the vector between the start \((p_i)\) and end \((p_j)\) points of the region (constraint \([b]\)). The first two constraints can be achieved using elementary translation and rotation of the dictionary coordinates. Thus transformed, the best rotation \( \theta^* \) of \( Q \) about the \( p_j - p_i \) axis is found so that constraint \([c]\) is realized; this constrained version of the problem can also be solved analytically using straightforward linear algebra.

Once the transformation of \( Q \) to \( Q' \) is achieved as described above, Alice uses it to transmit the interior points of the region, \( p_{i+1}, \ldots, p_{j-1}, \) by: (1) transmitting the best rotation about the \( p_j - p_i \) axis so that Bob, after receiving it, can orient the nominated dictionary fragment in the same orientation as Alice, and (2) transmitting the spatial deviations of the interior points with respect to the dictionary fragment.

Rotation \( \theta^* \) is transmitted over an uniform distribution on a circle whose radius \( r_{\theta^*} \) is the farthest distance of an interior point within the dictionary fragment from its axis of rotation. Dividing the circumference of a circle of radius \( r_{\theta^*} \) into arcs of length \( \epsilon \) and stating the arc number in which the rotated coordinate with the farthest radius to the axis falls gives the code length of stating \( \theta^* \):

\[
I(\theta^*) = -\log_2 \left( \frac{\epsilon}{2\pi r_{\theta^*}} \right) = \log_2(2\pi r_{\theta^*}) - \log_2 \epsilon \text{ bits.}
\]

Let the deviations of interior coordinates with respect to the corresponding (reoriented) dictionary fragment coordinates be denoted as \((\delta x_k, \delta y_k, \delta z_k)\). These components are transmitted using a normal distribution with a \( \mu \) of 0 and a standard deviation \( \sigma \) set to the sample standard deviation computed from these error components. (Note, \( \sigma \) is the root-mean-squared distance between corresponding coordinates after the aforementioned orthogonal transformation.)

The MML estimate to transmit a set of independent data \((\delta x_1, \delta y_1, \delta z_1), (\delta x_2, \delta y_2, \delta z_2), \ldots (\delta x_{|Q|-2}, \delta y_{|Q|-2}, \delta z_{|Q|-2})\) using a normal distribution is given by (see [12]):

\[
I(\delta x, \delta y, \delta z) = \frac{1}{2} (N - 1) \log_2 \sigma^2 + \frac{N - 1}{2} + \frac{N}{2} \log_2 \left( \frac{2\pi}{\epsilon^2} \right) + \frac{1}{2} \log_2 (2N^2) + \log_2 (R_\sigma) + 1 + \log_2 \kappa_1 \text{ bits.}
\]

where \( N = 3(|Q| - 2) \) is the total number of component spatial deviations being transmitted, \( R_\sigma \) gives the \((prior)\) range on \( \log_2 \sigma \), and \( \kappa_1 \approx \frac{1}{12} \) denotes the constant corresponding to quantizing lattices proposed by Conway and Sloane [17]. Here we bound \( \sigma \) at 3Å or less, because this is consistent with the empirically observed limits of utility of root-mean-squared-deviation measure to estimate similarity between rigid-body protein fragments.

Therefore, combining all the code lengths described so far, the code length required to transmit coordinates of a region of a protein using an assigned dictionary fragment is

\[
I_{\text{dict-frag}}(p_1, \ldots, p_j) = I_{\text{sphere}}(p_j) + I(\theta^*) + I(\delta x, \delta y, \delta z)
\]
5.2 Describing a protein region using the null model For encoding any region within a protein structure, there is an implicit null model which any dictionary fragment encoding has to beat in order to be chosen. This ensures that incompressible regions (or random coils) in proteins can be stated, without the support of any dictionary fragment. Using Equation 4.6, the code length to transmit a segment of a protein \( p_1, \ldots, p_j \) as a random coil takes:

\[
I_{\text{coil}}(p_1, \ldots, p_j) = I_{\text{null}}(p_1, \ldots, p_j) \quad \text{bits.}
\]

5.3 Describing a protein structure as a dissection of assigned dictionary fragments Here we describe the general language of communicating a protein \( \mathcal{P} \) given a dissection over a dictionary \( \mathcal{D} = \{ Q_1, Q_2, \ldots, Q_{|\mathcal{Q}|} \} \) of fragments. (The search for optimal dissection is dealt in Section 5.4.)

Formally, a dissection of \( \mathcal{P} = \{ p_1, \ldots, p_n \} \) gives an ordered subset of coordinates \( \mathcal{P}' = \{ p'_1 \equiv p_{i_1}, p'_2 \equiv p_{i_2}, \ldots, p'_m \equiv p_{i_m} \} \) where \( 1 = i_1 < i_2 < \cdots < i_m = n \). Each successive pair of coordinates in \( \mathcal{P}' \), \( \langle p'_1, p'_2 \rangle, \langle p'_2, p'_3 \rangle, \ldots, \langle p'_{m-1}, p'_m \rangle \), defines a pair of start and end points of a segment (or region). (Notice that the dissection \( \mathcal{P}' \) gives \( m-1 \) successive, non-overlapping regions of \( \mathcal{P} \), where end point of one segment is same as the start point of the next.) Assigned to each region \( \langle p'_k, p'_{k+1} \rangle \) containing \( l_k = i_{k+1} - i_k + 1 \) points is an integer \( t_k \) that specifies the model type assigned to that region in the protein. This integer ranges from \([1, |\mathcal{D}| + 1]\), accounting for possible assignments over any one of \(|\mathcal{Q}|\) available fragments plus the implicit random coil (null) model.

Given some dissection of \( \mathcal{P} \) using the dictionary \( \mathcal{D} \), Alice can transmit the coordinates over a two part message: First part of the message: In the first part of the message, Alice communicates the segmentation, and the corresponding assignments of models (dictionary fragment or coil) as a hypothesis on the observed coordinate data. To achieve this, the following information is transmitted. Using Equation 4.3 stating the number of segments \( (m-1) \) in the dissection takes \( I_{\log^*}(m-1) \) bits. The segmentation can be stated simply as a sequence of assigned model types. Assuming each model type is uniformly likely, each integer \( t_k \) corresponding to the assigned model type takes

\[
I_{\text{uniform}}(t_k) = \log_2(|\mathcal{D}| + 1) \quad \text{bits.}
\]

Thus, the message length of the first part requires

\[
I_{\text{first}}(\mathcal{P}'|\mathcal{D}) = I_{\log^*}(m-1) + \sum_{k=1}^{m-1} I_{\text{uniform}}(t_k) \quad \text{bits}
\]

Second part of the message: In the second part, Alice transmits the coordinates of \( \mathcal{P} \) using the hypothesis described in the first part of the message. The procedure to transmit coordinate data of a segment of a protein using dictionary fragment and coil models have already been discussed in sections 5.1 and 5.2.

The message length of the second part is given by:

\[
I_{\text{second}}(\mathcal{P}|\mathcal{P}'&\mathcal{D}) = \sum_{k=1}^{m-1} I_{\text{model}}(p'_k, \ldots, p'_{k+1}) \quad \text{bits}
\]

where \( I_{\text{model}}(p'_k, \ldots, p'_{k+1}) = I_{\text{dict-frag}}(p'_k, \ldots, p'_{k+1}) \) if \( t_k \) corresponds to a dictionary fragment and \( I_{\text{model}} = I_{\text{coil}}(p'_k, \ldots, p'_{k+1}) \) if \( t_k \) corresponds to a random coil.

Thus, the total message length of communicating the coordinates of a particular protein \( \mathcal{P} \) using some dissection \( \mathcal{P}' \) over a given dictionary \( \mathcal{D} \) is computed by combining the terms in Equations 5.12 and 5.13

\[
I(\mathcal{P}'&\mathcal{P}|\mathcal{D}) = I_{\text{first}}(\mathcal{P}'|\mathcal{D}) + I_{\text{second}}(\mathcal{P}|\mathcal{P}'&\mathcal{D})
\]
5.4 Search for the optimal dissection of a particular protein given a dictionary of substructures

Given Equation 5.14, the problem of inferring the best dissection can now be stated formally as follows: Given a protein \( P \) and a dictionary \( D \), find the dissection \( P' \) of \( P \) such that the message length to transmit \( P \) losslessly is minimized.

In the protein \( P \), any pair of points can potentially be the start and end points of a dissected segment. At the same time, any dissected segment can potentially be described using a coil model or any of the dictionary fragments (provided the length of the assigned dictionary element is same as the length of the segment). Therefore, the procedure to assign dictionary fragments to coordinates \( \{p_1, \cdots, p_{|P|}\} \) in a given protein begins by constructing a set of \(|D| + 1\) code length matrices, one for each possible encoding model type \( t \). Using Equations 5.12 and 5.13, the code length matrices are filled as:

\[
H^t(i, j) = I_{\log^*}(j - i) + I_{\text{uniform}}(t) + I_{\text{model}}(p_i, \cdots, p_j)
\]

where any cell \((i, j)\) of the matrix for type \( t \) gives the code length of stating the segment \( p_i \) to \( p_j \) using the model \( t \).

Optimal dissection as a dynamic program over the code length matrices, \( \{H^1, \cdots, H^{|D|}, H^{|D|+1}\} \): The dissection of \( P \) using various model types enforces a linear ordering constraint with optimal substructure and overlapping subproblems which can be composed incrementally. This allows the search for an optimal dissection of a protein structure, given a dictionary, to be formulated using a dynamic program that can be solved efficiently.

Let any \( M(i) \) store the optimal encoding message length to transmit points \( p_1, \cdots, p_i \), for all \( 1 \leq i \leq |P| \). With the boundary condition of \( M(1) = 0 \), the dynamic programming recurrence to find the optimal assignment is given by

\[
M(j) = \min_{i=1}^{j-1} \left\{ \begin{array}{l}
H^t(1, j) \\
M(i) + \min_t H^t(i, j)
\end{array} \right. \quad \text{model} = \begin{cases}
D \ni Q_t \\
t = \text{coil}
\end{cases}
\]

The above recurrence is used to fill the array \( M \) iteratively from 1 to \( |P| \). On completion, the best assignment of dictionary fragments can be derived by remembering the index \( i \) and type \( t \) from which the optimal \( M(j) \) is computed. The time complexity to find the optimal encoding is \( O(|D||P|^2) \).

6 Describing an entire collection of protein structures using a given dictionary of substructures

Restating Equation 3.1 in Section 3, \( I(D&C) = I(D) + I(C|D)) \). That is, the optimal encoding of a collection of protein structures \( C \) using a given dictionary \( D \) involves a one-off cost of stating the dictionary \( (I(D)) \), followed by the optimal encodings of the individual proteins in the collection \( (I(C|D)) \). Below we show the details of how the respective terms are estimated.

6.1 Estimating \( I(D) \) for a given dictionary For the dictionary \( D \) with substructures \( \{Q_1, Q_2, \cdots, Q_{|D|}\} \), the statement cost, \( I(D) \), involves: the transmission of the the number of substructures \( (|D|) \), followed by the statement cost of each fragment’s coordinates \( Q_t, D \) over a null model message.

From Equations 4.5 and 4.6 we can derive the message length of stating the dictionary \( D \) as:

\[
I(D) = I_{\log^*}(|D|) + \sum_{i=1}^{|D|} I_{\text{null}}(Q_t) \quad \text{bits}.
\]
6.2 Estimating $I(C|D)$ With the dictionary transmitted to Bob, Alice can now encode the structures in the collection using the method described in Section 5. Using Equation 5.14, the statement cost to transmit the coordinates of a collection of proteins $C = \{P_1, \cdots, P_{|C|}\}$ using the previously stated dictionary $D$ takes:

(6.17) 

$$I(C|D) = \sum_{i=1}^{|C|} I(P'_i & P_i | D)) \text{ bits}$$

7 The search for the optimal dictionary

Section 6.2 has established a rigorous objective to search for the dictionary of fragments building blocks of protein three-dimensional structures. It follows from the aforementioned objective (also Equation 5.1) that an optimal dictionary for a collection of protein structures is one for which the statement of the dictionary, plus the sum of the optimal encodings of all the proteins in the set, is the shortest. That is, the objective of this work is to find a $D^*$ such that:

(7.18) 

$$I(D^* & C) = \min_{\forall D} I(D & C) \text{ bits.}$$

Clearly, any fragment (of arbitrary size) from within any protein in the collection is a potential candidate for the dictionary. Therefore searching for the best dictionary leads to a very large optimization problem. Since the problem is computational intractable to find a provably optimal dictionary, we designed a simulated annealing algorithm in order to evolve a dictionary that iteratively converges to the best dictionary defined by the Equation 7.18.

Simulated Annealing is an heuristic approach which has an analogy with cooling of solids. Here, we consider each possible dictionary (of arbitrary number of fragments) as being analogous to some state of a physical system. The message length of transmitting a collection of structures using any dictionary given by Equation 5.1 is analogous to the internal energy of the physical system in that state.

The method starts with an empty dictionary. In this state, each protein in the collection is transmitted raw, as a random coil using the null model. The strategy involves iteratively perturbing the dictionary from this initial empty state to a state where the total message length objective is minimized. Perturbations At each step the current dictionary is perturbed randomly (which is akin to sampling some new nearby dictionary state). The choice of moving to the new state or remaining in the current one is decided probabilistically. Specifically, at each iteration, our method employs one of the following randomly chosen perturbations:

Add: Append to the current dictionary a new fragment. This fragment is from a randomly chosen protein from the collection, of random length.

Remove: Remove a randomly chosen fragment from the existing dictionary state.

Swap: Replace a randomly chosen fragment in the current dictionary with another randomly chosen fragment from the collection. This is equivalent to the sequence of perturbations: ‘Remove’ followed by an ‘Add’.

Perturb length: Expand or shrink a randomly chosen fragment from the current dictionary state by one residue, at the randomly chosen end. (Expanding a previously chosen fragment is achieved by remembering its locus – i.e., the fragment’s source protein structure and its offset in the source. This
Figure 1: Four fragments chosen from our dictionary our approach discovered. (a) $\frac{2}{3}$ turns of three-fold $\beta$-helix of length 31 residues. It occurs 26 times in our collection. See wwPDB 2IC7 for one such instance. (b) An exotic $\beta$-hairpin of length 29 residues. This occurs 20 times in the collection. See wwPDB 1UJU for one such instance. (c) A long $\beta$-hairpin of length 22 residues with 47 occurrences. See wwPDB 1JO8 for one such instance. (d) 1 turn of four-fold $\beta$-helix of length 21, which occurs 19 times in the collection.

allows any existing fragment to be elongated by an additional coordinate, based on the information available in its source structure.)

**Probability of acceptance or rejection of any perturbation.** Equation 3.1 gives the estimate of the negative logarithm of the joint probability of a dictionary and a collection. In Section 2 we have seen that the difference between the message lengths using any two different hypotheses (here, two dictionary states) gives the log-odds posterior ratio. This implies, if the total message length using a perturbed dictionary is $k$ bits shorter (conversely, longer) than the current state, then the perturbed dictionary is $2^k$ times more likely (conversely, unlikely) than the perturbed state.

The simulated annealing heuristic starts with a high (time-varying) parameter $t$ (after temperature). Let $I_{\text{current}}$ and $I_{\text{perturbed}}$ be the total message lengths using the current and perturbed states of a dictionary. During any iteration, if $\Delta I \equiv (I_{\text{current}} - I_{\text{perturbed}}) < 0$, the perturbed state is immediately accepted as the new current state, and the procedure continued. Otherwise, the perturbed state is accepted with a probability of $1/2^{\Delta I}$.  

**Cooling schedule.** For simulated annealing algorithms, the variation of the temperature parameter $T$ controls the evolution of the dictionary states. We set $T$ to a high 10,000 at the start. The parameter $T$ decays at a constant rate of 0.88. For each value of $t$ between 10,000 and 10, we carry out 10,000 perturbations, while for value of $T$ less than 10, we carry out 100,000 perturbations to the dictionary. The stopping criterion is when the number of iterations reaches 2 million iterations (which occurs at $t = 0.246$).

**Implementation:** A program to optimally encode a given collection of structures and then search for the best dictionary using simulated annealing was implemented in the C++ programming language. Message Passing Interface (MPI) was used to parallelize this program to run on a large computational cluster.

A MapReduce model was used to distribute the encoding tasks on the large cluster. In this model, a master node accepts a collection of protein structures as input and evenly divides them into smaller subsets of structures, where each subset is distributed to worker nodes. For every perturbation of a current dictionary state, each worker node computes the sum of message lengths to optimally encode each of the structures in its allocated subset. The worker node then returns the message length back to the master node, which collects all the message lengths and combines them to compute the total message length to encode all the structures in the given collection using the current state of the dictionary.
8 Results

The work resulted in a dictionary of 1711 fragments ranging in length from 4 to 31 amino acids. This dictionary was used to dissect the entire collection of 8992 source structures. The average root-mean-square (r.m.s.) deviation of orthogonal superposition of the dictionary fragments to the assigned regions is 0.29 Å. (Note, this does not include the separate statement and application of deviations which is part of the transmitted message, which would reduce 0.29 Å to 0.) The average, over all proteins in the set, of the maximum r.m.s. deviation of any model from all regions it encodes, is 1.23 Å. In Fig. 1 shows the visualization of four fragments chosen from the dictionary we discovered, of lengths 31, 29, 22 and 21 respectively. Previous methods, due to the length constraint (see Section 1), are unable to detect recurrent fragments this long.

Figure 2 shows the optimal dissection into dictionary fragments of the structure of flavodoxin from *Desulfovibrio vulgaris* (wwPDB entry 1J8Q, solved at 1.35 Å resolution [18].) To encode losslessly the flavodoxin structure, the dissection would be accompanied by the vector deviations of the Cα atoms of flavodoxin from the assigned canonical dictionary fragments. Noteworthy properties of the dissection in this example include:

1. The fits of the individual fragments of the dictionary to the structure are quite precise in almost all cases. The deviations are of the order of only tenths of an Å. (The maximum r.m.s.d. between the coordinates of a region in the protein and the assigned model is 0.53 Å; this occurs only once. All other r.m.s.d. values are ≤ 0.32 Å.)

2. Here the dictionary fragments account for the entire structure. No regions of the structure need be encoded as a random coil; that is, using the null model.

3. The range of lengths of the dictionary fragments appearing in the dissection of flavodoxin is from

| Resi | Model | RMSD (Å) |
|------|-------|----------|
| 2 - 6 | m1096 | 0.12     |
| 6 - 10 | m1195 | 0.21     |
| 10 - 13 | m1595 | 0.24     |
| 13 - 22 | m0231 | 0.16     |
| 22 - 32 | m0159 | 0.18     |
| 32 - 37 | m0874 | 0.27     |
| 37 - 41 | m1202 | 0.24     |
| 41 - 45 | m1246 | 0.18     |
| 45 - 49 | m1323 | 0.17     |
| 49 - 55 | m0502 | 0.19     |
| 55 - 60 | m0930 | 0.28     |
| 60 - 63 | m1685 | 0.07     |
| 63 - 67 | m1306 | 0.26     |
| 67 - 71 | m1194 | 0.16     |
| 71 - 78 | m0423 | 0.12     |

| Resi | Model | RMSD (Å) |
|------|-------|----------|
| 78 - 82 | m1415 | 0.14     |
| 82 - 85 | m1623 | 0.07     |
| 85 - 89 | m1083 | 0.20     |
| 89 - 92 | m1706 | 0.03     |
| 92 - 95 | m1611 | 0.06     |
| 95 - 100 | m0967 | 0.53     |
| 100 - 103 | m1571 | 0.17     |
| 103 - 115 | m0054 | 0.12     |
| 115 - 119 | m1499 | 0.19     |
| 119 - 123 | m1480 | 0.21     |
| 123 - 128 | m0769 | 0.29     |
| 128 - 133 | m0750 | 0.32     |
| 133 - 141 | m0281 | 0.12     |
| 141 - 148 | m0426 | 0.19     |

Figure 2: Dissection of the structure of flavodoxin from *Desulfovibrio vulgaris* (wwPDB entry 1J8Q). Parent structure shown in cyan; successive dictionary fragments in alternating red and green. Cofactor flavin mononucleotide (FMN) in ball-and-stick representation. Dissected regions are listed as a table below the picture.
Table 1: Clustering of models in the identified dictionary

| Class | Size | Code | Description                                      |
|-------|------|------|--------------------------------------------------|
| 1     | 616  | e    | short extended regions                           |
| 2     | 301  | t    | short non-hairpin turns (some $\beta$–bulges)    |
| 3     | 4    | t    | short non-hairpin turns                           |
| 4     | 325  | h    | short helices                                     |
| 5     | 164  | h    | short helices                                     |
| 6     | 167  | E    | extended regions (some with hooks at end)        |
| 7     | 6    | E    | extended regions (some curved)                   |
| 8     | 13   | T    | non-hairpin turns (some $\beta$–bulges)          |
| 9     | 50   | b    | shorter $\beta$–hairpin                          |
| 10    | 30   | B    | $\beta$–hairpins                                 |
| 11    | 3    | B    | $\beta$–hairpins, unequal arms (‘shepherd’s crook’) |
| 12    | 1    | B    | $\beta$–hairpins, unequal arms (‘shepherd’s crook’) |
| 13    | 3    | H    | irregular alpha helix (plus C$_\alpha$-only)     |
| 14    | 17   | H    | long alpha helices                                |
| 15    | 2    | $\Omega$ | long wide loops (Omega loop)                   |
| 16    | 1    | $\Omega$ | long wide loops (Omega loop)                   |
| 17    | 2    | $\Omega$ | long wide loops (Omega loop)                   |
| 18    | 1    | C    | double $\beta$–hairpin (‘paper clip’)            |
| 19    | 1    | B    | long twisted $\beta$–hairpins                   |
| 20    | 2    | $\Omega$ | helix-strand-helix-strand / wide (Omega) loops  |
| 21    | 1    | 3    | $\frac{1}{3}$ turns of three-fold $\beta$–helix |
| 22    | 1    | 4    | 1 turn of four-fold ($\beta$–helix)             |

4 to 13. Some of the segments correspond to individual fragments of secondary structure – helices and strands of sheets. Others correspond to N- or C-terminal parts of secondary structures, plus parts of the loops either preceding or succeeding them.

4. The sequence of dictionary fragments provides a one-dimensional representation of the protein folding pattern.

Clusteriing of the dictionary fragments: To further rationalize the 1711 dictionary fragments, we clustered the dictionary fragments into coarse structural classes with the UPGMA method [19] using the Mahalanobis distance [20] computed from the following characterising properties:
(1) the number of backbone hydrogen bonds between residues separated by 4 in the sequence (to group helices which demonstrate this periodicity),
(2) the distance between the C$_\alpha$ atoms of N- and C-terminal residues,
(3) the cosine of the angle between the C$_\alpha$ atom of the N-terminal residue, the C$_\alpha$ atom of the middle residue (or, for fragments containing even numbers of residues, the average position of the C$_\alpha$ atoms of...
the two middle residues),
(4) the r.m.s. deviation of a fit of the $C_\alpha$ atoms to a straight line, and
(5) the average value of the cosine of the dot products of $C\rightarrow O$ vectors of successive residues.

Table I in the main text shows the clusters, and suggested class codes.

Taking together the dissections of all proteins in the source set, the usage of different models in terms of case-insensitive coarse classes defined in Table I in main text. The distribution of usage, showing the prominence of secondary-structure elements, is:

| E   | H   | T   | B   | Ω   | C   | 3   | 4   |
|-----|-----|-----|-----|-----|-----|-----|-----|
| 180360 | 107262 | 72096 | 12390 | 447 | 48  | 28  | 21  |

The distribution of two-letter combinations is:

| EE  | HH  | HE  | HE  | EH  | 38735 |
|-----|-----|-----|-----|-----|-------|
| ET  | 38401 | TE  | 36721 | TH  | 21153 |
| TT  | 11583 | EB  | 9074 | BE  | 8797  |
| TB  | 1460  | HB  | 1247 | BH  | 1223  |
| EO  | 306   | OE  | 232  | OH  | 169   |
| HO  | 62    | EC  | 48   | CE  | 35    |
| 3E  | 22    | OB  | 18   | E4  | 15    |
| 4E  | 13    | CT  | 12   | T3  | 8     |
| BO  | 6     | 3T  | 3    | 33  | 3     |
| H3  | 1     | CB  | 1    |     |       |

Repetitions – EE and HH – lead the list, and that symmetric pairs (HE/EH, and ET/TE) tend to have similar distributions. The first 8 pairs account for 90% of the two-character combinations. This confirms the well-known importance of secondary structures as building blocks of proteins (assuming that most of the E regions are in fact parts of $\beta$–sheets; conversely, many of the $\beta$–hairpins include strands of sheet.) The first combination that does not contain any secondary structural element, TT ($T = \text{non-hairpin turn}$) accounts for 3% of the total number of pairwise combinations.

One-dimensional representation of protein folds. We consider as a case study the assignment of fragments to clusters which reveals structural patterns in Drosophila lebanonensis alcohol dehydrogenase (wwPDB code 1SBY).

Figure 3 shows the dissection of amino acid residues 1-183 of 1SBY into dictionary fragments. In addition to our dissection, we show the secondary structure assignments from the wwPDB file for the NAD-binding domain supersecondary structure.

The sequence of class symbols (column 6 in table in Figure 3), converted to upper case to suppress the distinction between short and long versions of the same substructure, affords a more perspicuous representation of this dissection:

```
EETTHTEEETHHEEETHHTTEETHHTEE
```

(E = strand, H = helix, T = non-hairpin turn (Table 1).)

This is an instance of the regular expression:

```
(E+T+H+T+H+T+T+E+H+E+H+T+T+T+H+T+T+E+) \{2\}
```

This sequence from the dissection provides a concise one-dimensional representation of the folding pattern. It captures the duplication of the two $\beta - \alpha - \beta - \alpha - \beta$ substructures (and the points of
| Model   | Residues | Length | r.m.s.d. | Class Code | Description of class       | Secondary Residues | Residues | Structure |
|---------|----------|--------|----------|------------|-----------------------------|--------------------|----------|-----------|
| m1685   | 1 - 4    | 4      | 0.168    | 1          | e short extended region     |                    |          |           |
| m0546   | 4 - 10   | 7      | 0.234    | 6          | E extended region           |                    |          |           |
| m1700   | 10 - 13  | 4      | 0.056    | 2          | t short turn                |                    |          |           |
| m1400   | 13 - 17  | 5      | 0.442    | 2          | t short turn                |                    |          |           |
| m0046   | 17 - 29  | 13     | 0.498    | 4          | h short helix               |                    |          |           |
| m1138   | 29 - 33  | 5      | 0.085    | 2          | t short turn                |                    |          |           |
| m0602   | 33 - 39  | 7      | 0.538    | 6          | E extended region           |                    |          |           |
| m1556   | 39 - 42  | 4      | 0.136    | 1          | e short extended region     |                    |          |           |
| m0195   | 42 - 52  | 11     | 0.174    | 4          | h short helix               |                    |          |           |
| m0967   | 52 - 57  | 6      | 0.183    | 1          | e short extended region     |                    |          |           |
| m0855   | 57 - 62  | 6      | 0.230    | 1          | e short extended region     |                    |          |           |
| m1247   | 62 - 66  | 5      | 0.170    | 2          | t short turn                |                    |          |           |
| m0243   | 66 - 74  | 9      | 0.200    | 4          | h short helix               |                    |          |           |
| m0316   | 74 - 82  | 9      | 0.142    | 4          | h short helix               |                    |          |           |
| m0791   | 82 - 87  | 6      | 0.304    | 1          | e short extended region     |                    |          |           |
| m1188   | 87 - 91  | 5      | 0.220    | 1          | e short extended region     |                    |          |           |
| m1291   | 91 - 95  | 5      | 0.299    | 1          | e short extended region     |                    |          |           |
| m0804   | 95 - 100 | 6      | 0.226    | 2          | t short turn                |                    |          |           |
| m0315   | 100 - 108| 9      | 0.130    | 4          | h short helix               |                    |          |           |
| m0096   | 108 - 119| 12     | 0.182    | 4          | h short helix               |                    |          |           |
| m1120   | 119 - 123| 5      | 0.168    | 2          | t short turn                |                    |          |           |
| m0600   | 123 - 129| 7      | 0.257    | 2          | t short turn                |                    |          |           |
| m0668   | 129 - 134| 6      | 0.413    | 1          | e short extended region     |                    |          |           |
| m1083   | 134 - 138| 5      | 0.215    | 1          | e short extended region     |                    |          |           |
| m1243   | 138 - 142| 5      | 0.140    | 5          | h short helix               |                    |          |           |
| m0408   | 142 - 149| 8      | 0.660    | 2          | t short turn                |                    |          |           |
| m0177   | 149 - 159| 11     | 0.160    | 4          | h short helix               |                    |          |           |
| m0180   | 150 - 169| 11     | 0.165    | 4          | h short helix               |                    |          |           |
| m1260   | 169 - 173| 5      | 0.151    | 2          | t short turn                |                    |          |           |
| m0636   | 173 - 179| 7      | 0.448    | 6          | E extended region           |                    |          |           |
| m1188   | 179 - 183| 5      | 0.152    | 1          | e short extended region     |                    |          |           |

Figure 3: Dissection of alcohol dehydrogenase structure (1SBY) using dictionary fragments.
insertion of non-pattern elements) but not their symmetrical spatial disposition. Nevertheless, the dissection provides a faithful signature of the NAD-binding domain folding pattern [21].

The rich repertoire of algorithms on character strings is applicable. For example, the string could be used to design regular expressions for probing collections for similar structures.

More generally, standard regular-expression-matching algorithms permit application of the linear representations to identify common folding patterns in a set of structures, or, specifically, to classify a newly-determined structure, as for example in SCOP [22] and CATH [23]. The representation can also identify variations and deviations from standard folding patterns in known families. For instance, some NAD-binding domains (including Drosophila lebanonensis alcohol dehydrogenase) contain extra helices and/or hairpins [24] and this would be revealed by a ‘sequence’ alignment of the dissections of members of this family.

9 Conclusion

This work introduces a novel method to infer the dictionary of building blocks of protein structures. This work fall squarely into one of the most important cross-disciplinary areas of modern science, where biology and computing meet. The approach described in this paper is a successful demonstration of rigorous statistical inference applied to an important data mining problem in structural Bioinformatics. The knowledge of this dictionary direct us to a number of avenues for further research. These include, for instance, straightforward generalisations, such as inclusion of backbone and even sidechain atoms. The exploration of the linear representation of folding patterns, and its potential correlation with sequence signals is a particularly attractive challenge. Working out the grammar associated with sequences of classes of dictionary fragments can illuminate different levels of folding architectures. The dictionary fragments themselves might well be applicable in approaches to predictions of protein structure, the holy grail of bioinformatics.

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