FINITE EXPRESSION METHOD FOR SOLVING HIGH-DIMENSIONAL PARTIAL DIFFERENTIAL EQUATIONS

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Abstract. Designing efficient and accurate numerical solvers for high-dimensional partial differential equations (PDEs) remains a challenging and important topic in computational science and engineering, mainly due to the “curse of dimensionality” in designing numerical schemes that scale in dimension. This paper introduces a new methodology that seeks an approximate PDE solution in the space of functions with finitely many analytic expressions and, hence, this methodology is named the finite expression method (FEX). It is proved in approximation theory that FEX can avoid the curse of dimensionality. As a proof of concept, a deep reinforcement learning method is proposed to implement FEX for various high-dimensional PDEs in different dimensions, achieving high and even machine accuracy with a memory complexity polynomial in dimension and an amenable time complexity. An approximate solution with finite analytic expressions also provides interpretable insights into the ground truth PDE solution, which can further help to advance the understanding of physical systems and design postprocessing techniques for a refined solution.

Key words. High-dimensional PDEs, Deep neural networks, Mathematical expressions, Finite expression method, Reinforcement learning.

1. Introduction. Partial differential equations (PDEs) play a fundamental role in scientific fields for modeling diverse physical phenomena, including diffusion [37, 27], fluid dynamics [2, 46] and quantum mechanics [16, 30]. Developing efficient and accurate solvers for numerical solutions to high-dimensional PDEs remains an important and challenging topic [49]. Many traditional solvers, such as finite element method (FEM) [40] and finite difference [18], are usually limited to low-dimensional domains since the computational cost increases exponentially in the dimension [49, 31]. Recently, neural networks (NNs) as mesh-free parameterization are widely employed in solving high-dimensional PDEs [13, 19, 25, 38, 47, 14] and control problems [21]. In theory, NNs have the capability of approximating various functions well and lessening the curse of dimensionality [45, 43, 51, 44, 24]. Yet the highly non-convex objective function and the spectral bias towards fitting a smooth function in NN optimization make it difficult to achieve high accuracy [50, 7, 41]. In practice, NN-based solvers can hardly achieve a highly accurate solution even when the true solution is a simple function, especially for high-dimensional problems [14, 33]. Besides, NN parametrization may still require large memory and high computation cost for high-dimensional problems [6]. Finally, numerical solutions from both traditional solvers and NN-based solvers lack interpretability, e.g., the dependence of the solution on variables is not readily apparent from the numerical solutions.

In this paper, we propose the finite expression method (FEX), a methodology that aims to find a solution in the function space of mathematical expressions with finitely many operators. Compared with the NN and FEM methods, our FEX enjoys the following advantages (summarized in Fig. 1.1): (1) The expression can reproduce the true solution and achieve high, even machine-level, accuracy. (2) The expression requires low memory (a line of string) for solution storage and low computational cost for evaluation on new points. (3) The expression has good interpretability with an explicit, readable form. Moreover, from an approximation theory perspective detailed in Sec. 2.3, the expression in FEX is capable of avoiding the curse of dimensionality in theory.

In FEX, we formulate the search for mathematical expressions as a combinatorial optimization (CO) involving both discrete and continuous variables. While many techniques can be used to solve the CO in FEX, we provide a numerically effective implementation based on reinforcement

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Fig. 1.1: Overview of various numerical solvers for PDEs. In our finite expression method, we aim to find a PDE solution as a mathematical expression with finitely many operators. The resulting solution can reproduce the true solution and achieve high, even machine-level, accuracy. Furthermore, the mathematical expression has memory complexity polynomial in dimension and provides interpretable insights to advance understanding of physical systems and aid in designing postprocessing techniques for solution refinement.

learning (RL). Traditional algorithms (e.g., genetic programming and simulated annealing [35]) address CO problems by employing hand-crafted heuristics that heavily rely on specific problem formulations and domain knowledge [11]. However, RL has emerged as a popular and versatile tool for learning how to construct CO solutions based on reward feedback, without the need for extensive heuristic design. The success of RL applications in CO, such as automatic algorithm design [39] and symbolic regression [36], has inspired us to seek mathematical expression solutions with RL. Specifically, in our implementation, the mathematical expression is represented by a binary tree, where each node is an operator along with parameters (scaling and bias) as shown in Fig. 3.1 and further explained in Sec. 3.1. The objective function we aim to minimize is a functional, and its minimization leads to the solution of the PDE, as described in Sec. 2.2. Consequently, our problem involves the minimization of both discrete operators and continuous parameters embedded within the tree structure. Optimizing both discrete and continuous variables simultaneously is inherently difficult. We propose a search loop for this CO as depicted in Fig. 1.2. Our idea is to first identify good operator selections that have a high possibility of recovering the true solution structure, and then optimize the parameters. The proposals of operator selections are drawn from a controller which will be updated via policy gradient [36] iteratively. In Sec. 4, we numerically demonstrate the ability of this RL-based implementation to find mathematical expressions that solve high-dimensional PDEs with high, even machine-level, accuracy. Furthermore, FEX provides interpretable insights into the ground truth PDE solutions, which can further advance understanding of physical systems and aid in designing postprocessing techniques to refine solutions.

2. An Abstract Methodology of Finite Expression Method. The goal of FEX is to find a mathematical expression to solve a PDE. This section will formally define the function space of mathematical expressions and formulate the problem of FEX as a CO.

2.1. Function Space with Finite Expressions in FEX. Mathematical expressions will be introduced to form the function space of FEX.
Fig. 1.2: Representation of the components of our FEX implementation. (a) The searching loop for the finite expression solution consists of expression generation, score computation, controller update, and candidate optimization. (b) Depiction of the expression generation with a binary tree and a controller $\chi$. The controller outputs probability mass functions for each node in the tree, and these functions are used to sample node values. The final expression, which incorporates learnable scaling and bias parameters, is constructed based on the predefined tree structure and the sampled node values.

**Definition 2.1 (Mathematical expression).** A mathematical expression is a combination of symbols, which is well-formed by syntax and rules and forms a valid function. The symbols include operands (variables and numbers), operators (e.g., “+”, “sin”, integral, derivative), brackets, punctuation.

In the definition of mathematical expressions, we only consider the expression that forms a valid function. In our context, “$\sin(x \times y) + 1$” is a mathematical expression, but, for example, “$5 > x$” and “$\sin(x \times y)$+” are not a mathematical expression as they do not form a valid function. The operands and operators comprise the structure of an expression. The parentheses play a role in clarifying the operation order.

**Definition 2.2 ($k$-finite expression).** A mathematical expression is called a $k$-finite expression if the number of operators in this expression is $k$.

“$\sin(x \times y) + 1$” is a 3-finite expression since there are three operators in it (“\times”, “sin”, and “+”). The series, such as “$1 + x^1 + \frac{x^2}{2} + \frac{x^3}{6} + \cdots$”, belongs to a mathematical expression, but it is not a finite expression since the amount of the operators is infinite. Formally, with the concept of finite expression, we can define FEX as follows,

**Definition 2.3 (Finite expression method).** The finite expression method is a methodology to solve a PDE numerically by seeking a finite expression such that the resulting function solves the PDE approximately.
We denote $S_k$ as a set of functions that are formed by finite expressions with the number of operators not larger than $k$. $S_k$ forms the function space of FEX. Clearly, $S_1 \subset S_2 \subset S_3 \cdots$.

2.2. Identifying PDE Solutions in FEX. We denote a functional $L : S \rightarrow \mathbb{R}$ associated with a given PDE, where $S$ is a function space and the minimizer of $L$ is the best solution to solve the PDE in $S$. In FEX, given the number of operators $k$, the problem of seeking a finite expression solution is formulated as a CO over $S_k$ via

\begin{equation}
(2.1) \quad \min \{L(u) | u \in S_k\}.
\end{equation}

The choice of the functional $L$ is problem-dependent and one may conceive a better functional for a specific PDE with a specific constraint or domain. Some popular choices include least-square methods \cite{28, 47, 12} and variation formulations \cite{14, 52, 8, 53}.

2.2.1. Least Square Method. Suppose that the PDE is given by

\begin{equation}
(2.2) \quad D u(x) = f(u(x), x), \quad x \in \Omega, \quad B u(x) = g(x), \quad x \in \partial \Omega,
\end{equation}

where $D$ is a differential operator, $f(u(x), x)$ can be a nonlinear function in $u$, $\Omega$ is a bounded domain in $\mathbb{R}^d$, and $B u = g$ characterizes the boundary condition (e.g., Dirichlet, Neumann and Robin \cite{15}). The least square method \cite{28, 47, 12} defines a straightforward functional to characterize the error of the estimated solution by

\begin{equation}
(2.3) \quad L(u) := \|D u(x) - f(u(x), x)\|_{L^2(\Omega)}^2 + \lambda \|B u(x) - g(x)\|_{L^2(\partial \Omega)}^2,
\end{equation}

where $\lambda$ is a positive coefficient to enforce the boundary constraint.

2.2.2. Variation Formulation. We next introduce the variation formulation, which is commonly used to identify numerical PDE solutions \cite{14, 52}. As an example, consider an elliptic PDE with homogeneous Dirichlet boundary conditions. This PDE is:

\begin{equation}
(2.4) \quad -\Delta u(x) + c(x) u(x) = f(x), \quad x \in \Omega, \quad u(x) = 0, \quad x \in \partial \Omega,
\end{equation}

where $c$ is a bounded function and $f \in L^2$. The solution $u$ to PDE (2.4) minimizes the variation formulation $\frac{1}{2} \int_{\Omega} \|\nabla u\|^2 + cu^2 \, dx - \int_{\Omega} f u \, dx$. By incorporating the boundary condition penalty into this variation, we obtain the functional:

\begin{equation}
(2.5) \quad L(u) := \frac{1}{2} \int_{\Omega} \|\nabla u\|^2 + cu^2 \, dx - \int_{\Omega} f u \, dx + \lambda \int_{\partial \Omega} u^2 \, dx.
\end{equation}

An alternative variation formulation can be defined using test functions. Let $v \in H^1_0(\Omega)$ be a test function, where $H^1_0(\Omega)$ denotes the Sobolev space whose weak derivative is $L^2$ integrable with zero boundary values. The weak solution $u$ of Eqn. (2.4) is defined as the function that satisfies the bilinear equations:

\begin{equation}
(2.6) \quad a(u, v) := \int_{\Omega} \nabla u \nabla v + cu v - f v \, dx = 0, \quad \forall v \in H^1_0(\Omega), \quad u(x) = 0, \quad x \in \partial \Omega,
\end{equation}

where $a(u, v)$ is constructed by multiplying (2.4) and $v$, and integration by parts. All derivatives of the solution function can be transferred to the test function through repeated integration by parts, yielding another bilinear forms \cite{8}. The weak solution can be reformulated as the solution to a saddle-point problem \cite{53}:

\begin{equation}
(2.7) \quad \min_{u \in H^1_0(\Omega)} \max_{v \in H^1_0(\Omega)} \frac{|a(u, v)|^2}{\|v\|^2_{L^2(\Omega)}}.
\end{equation}

Then, the functional $L$ identifying the PDE solution is:

\begin{equation}
(2.8) \quad L(u) := \max_{v \in H^1_0(\Omega)} \frac{|a(u, v)|^2}{\|v\|^2_{L^2(\Omega)}} + \lambda \int_{\partial \Omega} u^2 \, dx.
\end{equation}
2.3. Approximation Theory of Elementary Expressions in FEX. The most important theoretical question in solving high-dimensional problems is whether or not a solver suffers from the curse of dimensionality. It will be shown that the function space of $k$-finite expressions, i.e., $S_k$ in (2.1), is a powerful function space that avoids the curse of dimensionality in approximating high-dimensional continuous functions, leveraging the recent development of advanced approximation theory of deep neural networks [44, 24]. First of all, it can be proved that $S_k$ is dense in $C([0, 1]^d)$ for an arbitrary $d \in \mathbb{N}$ in the following theorem.

**Theorem 2.4.** The function space $S_k$, generated with operators including “+”, “-”, “$\times$”, “/”, “$\cdot | \cdot$”, “sign($\cdot$)”, and “$\lfloor \cdot \rfloor$”, is dense in $C([a, b]^d)$ for arbitrary $a, b \in \mathbb{R}$ and $d \in \mathbb{N}$ if $k \geq O(d^4)$.

Here “$\cdot | \cdot$”, “sign($\cdot$)”, and “$\lfloor \cdot \rfloor$” denote the absolute, sign and floor functions [44], respectively. The proof of Thm. 2.4 can be found in Appx. 6.1. The denseness of $S_k$ means that the function space of $k$-finite expressions can approximate any $d$-dimensional continuous functions to any accuracy, while $k$ is only required to be $O(d^4)$ independent of the approximation accuracy. The proof of Thm. 2.4 takes the advantage of operators “sign($\cdot$)” and “$\lfloor \cdot \rfloor$”, which might not be frequently used in mathematical expressions. If it is more practical to restrict the operator list to regular operators like “+”, “$\times$”, “$\sin(\cdot)$”, exponential functions, and the rectified linear unit (ReLU), then it can be proved that $S_k$ can approximate Hölder functions without the curse of dimensionality in the following theorem.

**Theorem 2.5.** Suppose the function space $S_k$ is generated with operators including “+”, “-”, “$\times$”, “$\div$”, “max{0, $x^\prime$}”, “$\sin(x)$”, and “$2^x$”. Let $p \in [1, +\infty)$. For any $f$ in the Hölder function class $\mathcal{H}_p^\alpha([0, 1]^d)$ and $\varepsilon > 0$, there exists a $k$-finite expression $\phi$ in $S_k$ such that $\|f - \phi\|_{L_p} \leq \varepsilon$, if $k \geq O(d^2(\log d + \log \frac{1}{\varepsilon})^2)$.

The proof of Thm. 2.5 can be found in Appx. 6.1. Although finite expressions have a powerful approximation capacity for high-dimensional functions, it is challenging to theoretically prove that our FEX solver to be introduced in Sec. 3 can identify the desired finite expressions with this power. However, the numerical experiments in Sec. 4 demonstrate that our FEX solver can identify the desired finite expressions to machine accuracy for several classes of high-dimensional PDEs.

3. An Implementation of FEX. Following the introduction of the abstract FEX methodology in Sec. 2, this section proposes a numerical implementation. First, binary trees are applied to construct finite expressions in Sec. 3.1. Next, our CO problem (2.1) is formulated in terms of parameter and operator selection to find expressions that approximate PDE solutions in Sec. 3.2. To resolve this CO, we propose implementing a search loop to identify effective operators that have the potential to recover the true solution when selected for expression. In Appendix 6.2, we provide pseudo-code for the FEX algorithm, which employs expanding trees to search for a solution. For the reader’s convenience, we have summarized the key notations used in this section in Table 3.1.

| Notation | Explanation |
|----------|-------------|
| $\mathcal{T}$ | A binary tree |
| $\epsilon$ | Operator sequence |
| $\theta$ | Trainable scaling and bias parameters |
| $\mathcal{L}$ | The functional associated with the PDE solution |
| $S$ | The scoring function that maps an operator sequence to [0, 1] |
| $\chi_\Phi$ | The controller parameterized by $\Phi$ |
| $J$ | The objective function for the policy-gradient approach |

Table 3.1: A summary of notations in the FEX implementation.
3.1. Finite Expressions with Binary Trees. A finite expression can be represented as a binary tree $T$, as depicted in Fig. 3.1. Each node in the tree is assigned a value from a set of operators, and all these node values collectively form an operator sequence $e$, following a predefined order to traverse the tree (e.g., inorder traversal).

Within each node featuring a unary operator, we incorporate two additional parameters, a scaling parameter $\alpha$ and a bias parameter $\beta$, to enhance expressiveness. All these parameters are denoted by $\theta$. Therefore, a finite expression can be denoted as $u(x; T, e, \theta)$, a function of $x$. For a fixed $T$, the maximal number of operators is upper bounded by a constant denoted $k_T$. In this implementation, $\{u(x; T, e, \theta) | e, \theta\}$ forms the function space in the CO to solve a PDE. This is a subset of functions expressible with at most $k_T$ finite expressions.

The configuration of the tree of various depths can be designed as in Fig. 3.1. Each tree node is either a binary operator or a unary operator that takes value from the corresponding binary or unary set. The binary set can be $\mathbb{B} := \{+, -, \times, \div, \cdots\}$. The unary set can be $\mathbb{U} := \{\sin, \exp, \log, \text{Id}, (\cdot)^2, \int \cdot dx_i, \frac{\partial}{\partial x_i}, \cdots\}$, which contains elementary functions (e.g., polynomial and trigonometric function), antiderivative and differentiation operators. Here “Id” denotes the identity map. Notice that if an integration or a derivative is used in the expression, the operator can be applied numerically. Each entry of the operator sequence $e$ corresponds one-to-one with traversing the tree nodes. The length of $e$ equals the total number of tree nodes. For example, when inorder traversal is used, Fig. 3.1 depicts a tree with 4 nodes and $e = (\text{Id}, \times, \exp, \sin)$.

The computation flow of the binary tree is conducted recursively. The operator of a node is granted higher precedence than that of its parent node. First, as in Fig. 3.1, we present the computation flow of the basic trees (a tree has a depth of 1 with only 1 operator). For a basic tree with a unary operator $u_1$, when the input is $i_0$, then the output $o_1 = \alpha u_1(i_0) + \beta$, where $\alpha$ and $\beta$ are scaling and bias parameters, respectively. For a basic tree with a binary operator $b_1$, when the input is $i_{01}$ and $i_{02}$, the output becomes $o_1 = b_1(i_{01}, i_{02})$. With these two basic trees, we are ready to define the computation for arbitrary depth by recursion, as the examples shown in Fig. 3.1. Specifically, the input of a parent node is the output of the child node(s).
vector to perform a linear transformation from $\mathbb{R}^d$ to $\mathbb{R}^1$.

3.2. Solving a CO in FEX. Given a tree $T$, we aim to seek the PDE solution from the function space $\{u(x; T, e, \theta) | e, \theta \} \subset \mathbb{S}_{k,T}$. The mathematical expression can be identified via the minimization of the functional $\mathcal{L}$ associated with a PDE, i.e.,

$$\min \{ \mathcal{L}(u(\cdot; T, e, \theta)) | e, \theta \}.$$  

We introduce the framework for implementing FEX, as displayed in Fig. 1.2a, to seek a minimizer of (3.1). The basic idea is to find a good operator sequence $e$ that may uncover the structure of the true solution, and then optimize the parameter $\theta$ to minimize the functional (3.1). In our framework, the searching loop consists of four parts: 1) Score computation (i.e., rewards in RL). A mix-order optimization algorithm is proposed to efficiently assess the score of the operator sequence $e$ to uncover the true structure. A higher score suggests a higher possibility to help to identify the true solution. 2) Operator sequence generation (i.e., taking actions in RL). A controller is proposed to generate operator sequences with high scores (see Fig. 1.2b). 3) Controller update (i.e., policy optimization in RL). The controller is updated to increase the probability of producing a good operator sequence via the score feedback of the generated ones. While the controller can be modeled in many ways (e.g., heuristic algorithm), we introduce the policy gradient in RL to optimize the controller. 4) Candidate optimization (i.e., a non-greedy strategy). During searching, we maintain a candidate pool to store the operator sequence with a high score. After searching, the parameters $\theta$ of high-score operator sequences are optimized to approximate the PDE solution.

3.2.1. Score Computation. The score of an operator sequence $e$ is critical to guide the controller toward generating good operator sequences and help to maintain a candidate pool of high scores. Intuitively, the score of $e$ is defined in the range $[0, 1]$, namely $S(e)$, by

$$S(e) := (1 + L(e))^{-1},$$  

where $L(e) := \min \{ \mathcal{L}(u(\cdot; T, e, \theta)) | \theta \}$. When $L(e)$ tends to 0, the expression represented by $e$ is close to the true solution, and the score $S(e)$ goes to 1. Otherwise, $S(e)$ goes to 0. The global minimizer of $\mathcal{L}(u(\cdot; T, e, \theta))$ over $\theta$ is difficult and expensive to obtain. Instead of exhaustively searching for a global minimizer, a first-order optimization algorithm and a second-order one are combined to accelerate the evaluation of $S(e)$.

First-order algorithms (e.g., the stochastic gradient descent [42] and Adam [26]) that utilize gradient to update are popular in machine learning. Typically, they demand a small learning rate and a substantial number of iterations to converge effectively. It can become time-consuming to optimize $L(e)$ using a first-order algorithm. Alternatively, second-order algorithms (e.g., the Newton method [3] and the Broyden-Fletcher-Goldfarb-Shanno method (BFGS) [17]) use the (approximated) Hessian matrix for faster convergence, but obtaining a good minimizer requires a good initial guess. To expedite the optimization process of $L(e)$ in our implementation, we employ a two-step approach. Initially, a first-order algorithm is utilized for $T_1$ steps to obtain a well-informed initial guess. Subsequently, a second-order algorithm is applied for an additional $T_2$ steps to further refine the solution. Let $\theta_0^e$ be an initialization and $\theta_{T_1+T_2}^e$ be the parameter set after $T_1 + T_2$ steps of this two-step optimization. Then $\theta_{T_1+T_2}^e$ serves as an approximate of $\arg \min_{\theta} \mathcal{L}(u(\cdot; T, e, \theta))$. Finally, $S(e)$ is estimated by

$$S(e) \approx (1 + \mathcal{L}(u(\cdot; T, e, \theta_{T_1+T_2}^e)))^{-1}.$$  

Remark that the approximation may exhibit significant variation due to the randomness associated with the initialization of $\theta_0^e$. 


3.2.2. Operator Sequence Generation. The role of the controller is to generate operator sequences with high scores during the searching loop. Let $\chi_\Phi$ be a controller with model parameter $\Phi$, and $\Phi$ is updated to increase the probability for good operator sequences during the searching loop. We use $e \sim \chi_\Phi$ to denote the process to sample an $e$ according to the controller $\chi_\Phi$.

Treating tree node values of $\mathcal{T}$ as random variables, the controller $\chi_\Phi$ outputs probability mass functions $p^1_\Phi, p^2_\Phi, \ldots, p^s_\Phi$ to characterize their distributions, where $s$ is the total number of nodes. Each tree node value $e_j$ is sampled from $p^i_\Phi$ to obtain an operator. Then $e := (e_1, e_2, \ldots, e_s)$ is the operator sequence sampled from $\chi_\Phi$. See Fig. 1.2 for an illustration. Besides, we adopt the $\epsilon$-greedy strategy [48] to enhance exploration of a potentially high-score $e$. With probability $\epsilon < 1$, $e_i$ is sampled from a uniform distribution of the operator set. With probability $1 - \epsilon$, $e_i \sim p^i_\Phi$. A larger $\epsilon$ leads to a higher probability to explore new sequences.

3.2.3. Controller Update. The goal of the controller update is to guide the controller toward generating high-score operator sequences $e$. The updating rule of a controller can be designed based on heuristics (e.g., genetic and simulated annealing algorithms) and gradient-based methods (e.g., policy gradient and darts [32]). As proof of concept, we introduce a policy-gradient-based updating rule in RL. The policy gradient method aims to maximize the return by optimizing a parameterized policy and the controller in our problem plays the role of a policy.

In this paper, the controller $\chi_\Phi$ is modeled as a neural network parameterized by $\Phi$. The training objective of the controller is to maximize the expected score of a sampled $e$, i.e.,

$$J(\Phi) := \mathbb{E}_{e \sim \chi_\Phi} S(e).$$

Taking the derivative of (3.4) with respect to $\Phi$, we have

$$\nabla_\Phi J(\Phi) = \mathbb{E}_{e \sim \chi_\Phi} \{S(e) \sum_{i=1}^{s} \nabla_\Phi \log(p^i_\Phi(e_i))\},$$

where $p^i_\Phi(e_i)$ is the probability corresponding to the sampled $e_i$. When the batch size is $N$ and $\{e^{(1)}, e^{(2)}, \ldots, e^{(N)}\}$ are sampled under $\chi_\Phi$ each time, the expectation can be approximated by

$$\nabla_\Phi J(\Phi) \approx \frac{1}{N} \sum_{k=1}^{N} \{S(e^{(k)}) \sum_{i=1}^{s} \nabla_\Phi \log(p^i_\Phi(e_i^{(k)}))\}.$$ 

Next, the model parameter $\Phi$ is updated via the gradient ascent with a learning rate $\eta$, i.e.,

$$\Phi \leftarrow \Phi + \eta \nabla_\Phi J(\Phi).$$

The objective in (3.4) helps to improve the average score of generated sequences. In our problem, the goal is to find $e$ with the best score. To increase the probability of obtaining the best case, the objective function proposed in [36] is applied to seek the optimal solution via

$$J(\Phi) = \mathbb{E}_{e \sim \chi_\Phi} \{S(e) | S(e) \geq S_{\nu, \Phi}\},$$

where $S_{\nu, \Phi}$ represents the $(1 - \nu) \times 100\%$-quantile of the score distribution generated by $\chi_\Phi$. In a discrete form, the gradient computation becomes

$$\nabla_\Phi J(\Phi) \approx \frac{1}{N} \sum_{k=1}^{N} \{(S(e^{(k)}) - \hat{S}_{\nu, \Phi}) \mathbbm{1}_{\{S(e^{(k)}) \geq \hat{S}_{\nu, \Phi}\}} \sum_{i=1}^{s} \nabla_\Phi \log(p^i_\Phi(e_i^{(k)}))\},$$

where $\mathbbm{1}$ is an indicator function that takes value 1 if the condition is true otherwise 0, and $\hat{S}_{\nu, \Phi}$ is the $(1 - \nu)$-quantile of the scores $\{S(e^{(i)})\}_{i=1}^{N}$. 

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3.2.4. Candidate Optimization. As introduced in Sec. 3.2.1 the score of \( e \) is based on the optimization of a nonconvex function at a random initialization. Therefore, the optimization may get stuck at poor local minimizers and the score sometimes may not reflect whether \( e \) reveals the structure of the true solution. The operator sequence \( e \) corresponding to the true solution (or approximately) may not have the best score. For the purpose of not missing good operator sequences, a candidate pool \( \mathbb{P} \) with capacity \( K \) is maintained to store several \( e \)’s of a high score.

During the search loop, if the size of \( \mathbb{P} \) is less than the capacity \( K \), \( e \) will be put in \( \mathbb{P} \). If the size of \( \mathbb{P} \) reaches \( K \) and \( S(e) \) is larger than the smallest score in \( \mathbb{P} \), then \( e \) will be appended to \( \mathbb{P} \) and the one with the least score will be removed. After the searching loop, for every \( e \in \mathbb{P} \), the objective function \( L(u(\cdot; T, e, \theta)) \) is optimized over \( \theta \) using a first-order algorithm with a small learning rate for \( T \) iterations.

4. Numerical Examples. Numerical results will be provided to demonstrate the effectiveness of our FEX implementation introduced in Sec. 3.2 using two classical PDE problems: high-dimensional PDEs with constraints (such as Dirichlet boundary conditions and integration constraints) and eigenvalue problems. The computational tools for high-dimensional problems are very limited and NNs are probably the most popular ones. Therefore, FEX will be compared with NN-based solvers. Through our examples, the goal is to numerically demonstrate that:

- **Accuracy.** FEX can achieve high and even machine accuracy for high-dimensional problems, while NN-based solvers can only achieve the accuracy of \( \mathcal{O}(10^{-4}) \) to \( \mathcal{O}(10^{-2}) \). Furthermore, we demonstrate the effectiveness of our RL-based approach by comparing its performance to a solution developed using genetic programming (GP) [35].

- **Scalability.** FEX is scalable in the problem dimension with an almost constant accuracy and a low memory requirement, i.e., the accuracy of FEX remains essentially the same when the dimension grows, while NN-based solvers have a worse accuracy when the dimension becomes larger.

- **Interpretability.** FEX provides interpretable insights of the ground truth PDE solution and helps to design postprocessing techniques for a refined solution.

In particular, to show the benefit of interpretability, we will provide examples to show that the explicit formulas of FEX solutions help to design better NN-parametrization in NN-based solvers to achieve higher accuracy. The FEX-aided NN-based solvers are referred to as FEX-NN in this paper. Finally, we will show the convergence of FEX with the growth of the tree size when the true solution can not be exactly reproduced by finite expressions using the available operators and a binary tree. All results of this section are obtained with 6 independent experiments to achieve their statistics.

4.1. Experimental Setting. This part provides the setting of FEX and NN-based solvers. The accuracy of a numerical solution \( \hat{u} \) compared with the true solution \( u \) is measured by a relative \( L^2 \) error, i.e., \( \| \hat{u} - u \|_{L^2(\Omega)} / \| u \|_{L^2(\Omega)} \). The integral in the \( L^2 \) norm is estimated by the Monte Carlo integral for high-dimensional problems.

**Implements of FEX.** The depth-3 binary tree (Fig. 1.2b) with 3 unary operators and 1 binary operator is used to generate mathematical expressions. The binary set is \( \mathbb{B} = \{ +, -, \times \} \) and the unary set is \( \mathbb{U} = \{ 0, 1, \text{Id}, (\cdot)^2, (\cdot)^3, (\cdot)^4, \exp, \sin, \cos \} \). A fully connected NN is used as a controller \( \mathbf{X}_\theta \) with constant input (see Appendix 6.3 for more details). The output size of the controller NN is \( n_1 |\mathbb{B}| + n_2 |\mathbb{U}| \), where \( n_1 = 1 \) and \( n_2 = 3 \) represent the number of binary and unary operators, respectively, and \( | \cdot | \) denotes the cardinality of a set.

There are four main parts in the implementation of FEX as introduced in Sec. 3.2. We will only briefly describe the key numerical choices here. (1) **Score computation.** The score is updated first by Adam with a learning rate 0.001 for \( T_1 = 20 \) iterations and then by BFGS with a learning rate 1 for
maximum $T_2 = 20$ iterations. (2) **Operator sequence generation.** The depth-3 binary tree (Fig. 1.2b) with 3 unary operators and 1 binary operator is used to generate mathematical expressions. The binary set is $\mathbb{B} = \{+, -, \times\}$ and the unary set is $\mathbb{U} = \{0, 1, \text{Id}, (\cdot)^2, (\cdot)^3, (\cdot)^4, \exp, \sin, \cos\}$. A fully connected NN is used as a controller $\mathbf{x}_\theta$ with constant input. The output size of the controller NN is $n_1|\mathbb{B}| + n_2|\mathbb{U}|$, where $n_1 = 1$ and $n_2 = 3$ represent the number of binary and unary operators, respectively, and $|\cdot|$ denotes the cardinality of a set. (3) **Controller update.** The batch size for the policy gradient update is $N = 10$ and the controller is trained for 1000 iterations using Adam with a fixed learning rate 0.002. Especially in Section of “Numerical Convergence”, since the deeper trees are used, the controller is updated for 5000 iterations for trees with different depth. We adopt the $\epsilon$-greedy strategy to increase the exploration of new $e_i$. The probability $\epsilon$ of sampling an $e_i$ by random is 0.1. (4) **Candidate optimization.** The candidate pool capacity is set to be $K = 10$. For any $e \in P$, the parameter $\theta$ is optimized using Adam with an initial learning rate 0.01 for $T_3 = 20,000$ iterations. The learning rate decays following a cosine decay schedule [23].

**Implements of NN-based Solvers.** Residual networks (ResNets) [22] parameterized by $\Theta$ are used to approximate a solution and a minimization problem $\min_{\Theta} \mathcal{L}(u(\cdot; \Theta))$ is solved to identify a numerical solution [47, 49, 14].

The ResNet maps from $\mathbb{R}^d$ to $\mathbb{R}^1$ and consists of seven fully connected layers with three skip connections. Each hidden layer contains 50 neurons. The neural network is optimized using the Adam optimizer with an initial learning rate 0.001 for 15,000 iterations. The learning rate is decayed following a cosine decay schedule.

**Poisson equation.** The coefficient $\lambda$ in the functional [2, 3] is 100. The batch size for the interior and boundary is 5,000 and 1,000, respectively. In the NN method, we use the ReLU$^2$ activation, i.e., $(\max\{x, 0\})^2$, in ResNet to approximate the solution.

**Linear conservative law.** The coefficient $\lambda$ in the functional [4, 3] is 100. In the NN method, we use ReLU $(\max\{x, 0\})$ activation. The batch size for the interior and boundary is 5,000 and 1,000, respectively. We use the same batch size as the NN method except that we increase the batch size to 20,000 for the interior and 4,000 for the boundary when the dimension is not smaller than 36.

**Schrödinger equation.** The coefficient $\lambda$ in the functional [4, 6] is 1. The batch size for estimating the first term and second term of [4, 6] is 2,000 and 10,000, respectively. Besides, ReLU$^2$ is used in ResNet.

**Implements of GP.** GP is an evolutionary algorithm used for automated symbolic regression and the evolution of computer programs to solve complex problems. Through a process of selection, crossover, and mutation, GP evolves and refines these programs over generations to optimize a specific fitness or objective function. We utilize GP to optimize the CO [2, 1] directly using the Python package GPlearn [1].

**4.2. High-dimensional PDEs.** Several numerical examples for high-dimensional PDEs including linear and nonlinear cases are provided here. In these tests, the true solutions of PDEs have explicit formulas that can be reproduced by the binary tree defined in Sec. 3.1 and $\mathbb{B}, \mathbb{U}$ defined in Sec. 4.1.

**4.2.1. Poisson Equation.** We consider a Poisson equation [14] with a Dirichlet boundary condition on a $d$-dimensional domain $\Omega = [-1, 1]^d$,

$$-\Delta u = f \text{ for } x \in \Omega, \quad u = g \text{ for } x \in \partial \Omega. \quad (4.1)$$

Let the true solution be $\frac{1}{2} \sum_{i=1}^d x_i^2$, and then $f$ becomes a constant function $-d$. The functional $\mathcal{L}$ is defined by least-square error [2, 3] is applied in the NN-based solver and FEX to seek the PDE solution for various dimensions ($d = 10, 20, 30, 40$ and 50).
4.2.2. Linear Conservation Law. The linear conservation law [9] is considered here with a domain $T \times \Omega = [0, 1] \times [-1, 1]^d$ and an initial value problem

\begin{equation}
\frac{\pi d}{4} u_t - \sum_{i=1}^{d} u_{x_i} = 0 \quad \text{for } x = (x_1, \cdots, x_d) \in \Omega, t \in [0, 1], \quad u(0, x) = \sin\left(\frac{\pi}{4} \sum_{i=1}^{d} x_i\right) \quad \text{for } x \in \Omega,
\end{equation}

where the true solution is $u(t, x) = \sin(t + \frac{\pi}{4} \sum_{i=1}^{d} x_i)$, and $d = 5, 11, 17, 23, 29, 35, 41, 47$ and $53$ in our tests. The functional $L$ in the NN-based solver and FEX to identify the solution is defined by

\begin{equation}
L(u) := \left\| \frac{\pi d}{4} u_t - \sum_{i=1}^{d} u_{x_i} \right\|_{L^2(T \times \Omega)}^2 + \lambda \left\| u(0, x) - \sin\left(\frac{\pi}{4} \sum_{i=1}^{d} x_i\right) \right\|_{L^2(\Omega)}^2.
\end{equation}

4.2.3. Nonlinear Schrödinger Equation. We consider a nonlinear Schrödinger equation [20] with a cubic term on a $d$-dimensional domain $\Omega = [-1, 1]^d$,

\begin{equation}
-\Delta u + u^3 + Vu = 0 \quad \text{for } x \in \Omega,
\end{equation}

where $V(x) = -\frac{1}{9} \exp\left(\frac{2}{d} \sum_{i=1}^{d} \cos x_i\right) + \sum_{i=1}^{d} \left(\frac{\sin^2 x_i}{d} - \frac{\cos x_i}{d}\right)$ for $x = (x_1, \cdots, x_d)$. And we let $\hat{u}(x) = \exp\left(\frac{1}{2} \sum_{j=1}^{d} \cos(x_j)/3\right)$ be the solution of the PDE (4.4). To avoid the trivial zero solution, we apply different strategies during the score computation and candidate optimization phases. During the score computation phase, the norm of the test function $u$, i.e., $\|u\|_{L^2(\Omega)}$, is used as a penalty for the function that is close to zero by

\begin{equation}
L_1(u) := \| -\Delta u + u^3 + Vu \|_{L^2(\Omega)}^2/\|u\|_{L^2(\Omega)}^3.
\end{equation}

During the candidate optimization phase, an integration constraint is imposed to (4.4), i.e., $\int_{\Omega} u(x) dx = \int_{\Omega} \hat{u}(x) dx$. The functional $L$ used in FEX to fine-tune the identified operator sequence is defined by

\begin{equation}
L_2(u) := \left\| -\Delta u + u^3 + Vu \right\|_{L^2(\Omega)}^2 + \lambda \left( \int_{\Omega} u(x) dx - \int_{\Omega} \hat{u}(x) dx \right)^2,
\end{equation}

where the second term imposes the integration constraint. The functional (4.6) is also used in the NN-based solver to approximate the PDE solution. Various dimensions are tested in the numerical results, e.g., $d = 6, 12, 18, 24, 30, 36, 42$ and $48$. Remark that we avoid using (4.6) in the score computation because the Monte-Carlo error tends to be significant in the second term of (4.6) when the batch size is small, but using a large batch size can lead to inefficient computations. Hence, for the computational efficiency, we will utilize (4.5) instead in the score computation, rather than (4.6).

4.3. Results. Three main sets of numerical results for the PDE problems above will be presented. First, the errors of the numerical solutions by NN-based solvers and FEX are compared. Second, a convergence test is analyzed when the tree size of FEX increases. Finally, FEX is applied to design special NN parametrization to solve PDEs in NN-based solvers.

Estimated Solution Error. The depth-3 binary tree (Fig. 3.1) is used in FEX with four nodes (a root node (R), a middle node (M), and two leave nodes (L1 and L2)). Fig. 4.1 shows the operator distribution obtained by FEX and the error comparison between the NN method and our FEX. The results show that NN solutions have numerical errors between $O(10^{-4})$ and $O(10^{-2})$ and the errors grow in the problem dimension $d$, agreeing with the numerical observation in the literature. Meanwhile, FEX can identify the true solution structure for the Poisson equation and
the linear conservation law with errors of order $10^{-7}$, reaching the machine accuracy since the single-float precision is used. In the results of the nonlinear Schrödinger equation, FEX identifies the solutions of the form $\exp(\cos(\cdot))$ but achieves errors of order $10^{-5}$. Note that $\int_{\Omega} \hat{u}(\mathbf{x})d\mathbf{x}$ in (4.6) is estimated by the Monte-Carlo integration with millions of points as an accurate and precomputed constant, but $\int_{\Omega} u(\mathbf{x})d\mathbf{x}$ can only be estimated with fixed and small batch size, typically less than 10,000, in the optimization iterations. As the dimension grows, the estimation error of $\int_{\Omega} u(\mathbf{x})d\mathbf{x}$ increases, and, hence, even the ground true solution has an increasingly large error according to (4.6). Therefore, the optimization solver may return an approximate solution without machine accuracy. Designing a functional $\mathcal{L}$ free of the Monte-Carlo error (e.g., Eqns. (2.3) and (4.3)) for the nonlinear Schrödinger equation could ensure machine accuracy. Furthermore, as shown in Figure 4.1, GP tends to identify solutions of lower quality that fail to achieve a high level of accuracy.

**Numerical Convergence.** The numerical convergence analysis is performed using the Poisson equation as an example. Binary trees of depths 2, 3, 4, 6 and 8 are used (see Fig. 3.1). The square operator $(\cdot)^2$ is excluded in $\mathbb{U}$ so that the binary tree defined in Sec. 3.1 can not reproduce the true solution (sum of the square of coordinates) exactly. This setting can mimic the case of a complicated solution while a small binary tree was used. Fig. 4.2 shows the error distribution of FEX with the growth of dimensions and the change of tree depths. FEX obtains smaller errors with increasing tree size. Notice that, compared with the errors of NN-based solvers reported in Fig. 4.1, FEX gets a higher accuracy when a larger tree is used.

**FEX-NN.** FEX provides interpretable insights of the ground truth PDE solution by the operator distribution obtained from the searching loop. It may be beneficial to design NN models with a special structure to increase the accuracy of NN-based solvers. In the results of the Poisson equation, we observe that the square operator has a high probability to appear at the leave nodes, which suggests that the true solution may contain the structure $\mathbf{x}^2 := (x_1^2, \ldots, x_d^2)$ at the input $\mathbf{x}$. As a result, we define the FEX-NN by $v(\mathbf{x}^2; \Theta)$ for the Poisson equation. Similarly, we use FEX-NNs $\sin(v(\mathbf{x}^2; \Theta))$ for the linear conservation law and $\exp(v(\mathbf{x}^2; \Theta))$ for the nonlinear Schrödinger equation. Fig. 4.1 shows the errors of FEX-NN with the growth of dimensions, and it is clear that FEX-NN outperforms the vanilla NN-based method by a significant margin.

**4.4. Eigenvalue Problem.** Consider identifying the smallest eigenvalue $\gamma$ and the associated eigenfunction $u$ of the eigenvalue problem [14].

(4.7) \[-\Delta u + wu = \gamma u, \quad \mathbf{x} \in \Omega, \quad \text{for } u = 0 \text{ on } \partial \Omega.\]

The minimization of the Rayleigh quotient $\mathcal{I}(u) = \frac{\int_{\Omega} \|
abla u\|^2d\mathbf{x} + \int_{\Omega} wu^2d\mathbf{x}}{\int_{\Omega} u^2d\mathbf{x}}$, s.t., $u|_{\partial \Omega} = 0$, gives the smallest eigenvalue and the corresponding eigenfunction. In the NN-based solver [14], the following functional is defined

(4.8) \[\mathcal{L}(u) := \mathcal{I}(u) + \lambda_1 \int_{\Omega} u^2d\mathbf{x} + \lambda_2 \left( \int_{\Omega} u^2d\mathbf{x} - 1 \right)^2\]

to seek an NN solution. Considering an example of $w = \|\mathbf{x}\|^2$ and $\Omega = \mathbb{R}^d$, the smallest eigenvalue of (4.7) is $d$ and the associated eigenfunction is $\exp(-\|\mathbf{x}\|^2/d)$. The domain $\Omega$ is truncated from $\mathbb{R}^d$ to $[-3, 3]^d$ for simplification as done in [14].

In FEX, the functional (4.8) is also used to estimate a solution. FEX discovers a high probability to have the “exp” operator at the tree root (100% for $d = 2, 4, 6, 8$, and $93.3\%$ for $d = 10$ as shown in Fig. 4.3). Therefore, it is reasonable to assume that the eigenfunction is of the form $\exp(v(\mathbf{x}))$. 

Let \( u(\mathbf{x}) = \exp(v(\mathbf{x})) \) and then Eqn. (4.7) is simplified to

\[
-\Delta v - \|\nabla v\|_2^2 + \|\mathbf{x}\|_2^2 = \gamma.
\]

Eqn. (4.9) does not have a trivial zero solution so we can avoid the integration constraint used in Eqns. (4.6) and (4.8), which leads to Monte-Carlo errors. Using Equation (4.9) and the Rayleigh quotient \( \mathcal{I} \), the values of \( v \) and \( \gamma \) are alternatively updated until they reach convergence. The detail of this iterative algorithm is presented in Appx. 6.4. Fig. 4.3 shows the relative absolute error of the estimated eigenvalues with the growth of the dimensions. We can see that directly optimizing (4.8) with the NN method produces a large error on the eigenvalue estimation, especially when the dimension is high (e.g., the relative error is up to 20% when \( d = 10 \)). With the postprocessing algorithm with FEX, we can identify the eigenvalue with an error close to zero.

5. Discussion and perspectives. In this paper, we proposed the finite expression method - a methodology to find PDE solutions as simple mathematical expressions. Our theory showed that mathematical expressions can overcome the curse of dimensionality. We provided one implementation of representing mathematical expressions using trees and solving the formulated combinatorial optimization in FEM using reinforcement learning. Our results demonstrated effectiveness of FEM at achieving high, even machine-level accuracy on various high-dimensional PDEs while existing solvers suffered from low accuracy in comparison.

While CO is inherently a complex challenge, continued exploration can be highly advantageous in achieving improved performance, particularly when dealing with more intricate PDE problems within the framework of FEX.

- Efficient computation of operator sequence score. We employed Formulation (3.3) as a proxy for assessing the quality of an operator sequence. However, it’s important to note that this formulation necessitates optimization over multiple steps, resulting in non-negligible computational costs that can impact the overall expense of solving the CO problem in FEX. Therefore, it is crucial to define a more efficient scoring method that reduces computational expenses and simplifies the identification of favorable operator sequences.

- Design of the controller. As an illustrative example of CO solving in this work, we employed a straightforward fully connected network to model the distribution responsible for proposing operator selections. It’s worth noting that more sophisticated techniques can also be applied, such as recurrent neural networks [36], which take prior decisions as input and generate new decisions as output. These advanced methods can offer enhanced modeling capabilities and adaptability to the context of the problem.

Data and Code Availability. No date is generated in this work. Source codes for reproducing the results in this paper are available online at: https://github.com/LeungSamWai/Finite-expression-method. The source codes are released under MIT license.

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6. Appendix.

6.1. Proofs of Theorems. Proof. [Proof of Thm. 2.4] By Thm. 1.1 of Shen et al (2021) [11], for any \( f \in C([a, b]^d) \) as a continuous function on \([a, b]^d\), there exists a fully connected neural network (FNN) \( \phi \) with width \( N = 36d(2d + 1) \) and depth \( L = 11 \) (i.e., 11 hidden layers) such that, for an arbitrary \( \varepsilon > 0 \), \( \| \phi - f \|_{L^\infty([a, b]^d)} < \varepsilon \). This FNN is constructed via an activation function with an explicit formula \( \sigma(x) = \sigma_1(x) := |x - 2|x^2 - 1| \) for \( x \in [0, \infty) \) and \( \sigma(x) = \sigma_2(x) := \frac{x}{|x| + 1} \) for \( x \in (-\infty, 0) \). Therefore, \( \sigma(x) = \frac{\sin(x) + 1}{2}\sigma_1(x) - \frac{\sin(x) - 1}{2}\sigma_2(x) \). Hence, it requires at most 18 operators to evaluate \( \sigma(x) \). For an FNN of width \( N \) and depth \( L \), there are \( N(d + 1) + (L - 1)N^2 \) operators “\( \times \)”, \( Nd - 1 + (L - 1)N(N - 1) \) operators “\(+\)”, and \( NL \) evaluations of \( \sigma(x) \) to evaluate an output of the FNN. Therefore, the FNN \( \phi \) is a mathematical expression with at most \( k_d := 103680d^4 + 103824d^3 + 39600d^2 + 6804d - 1 = \mathcal{O}(d^4) \) operators. Therefore, for any \( \varepsilon > 0 \), any continuous function \( f \) on \([a, b]^d\), there is a \( k_d \)-finite expression that can approximate \( f \) uniformly well on \([a, b]^d\) within \( \varepsilon \) accuracy. Since \( k_d \) is independent of \( \varepsilon \), it is clear that the function space of \( k_d \)-finite expressions is dense in \( C([a, b]^d) \). \( \square 

Proof. [Proof of Thm. 2.5] By Cor. 3.8 of Jiao et al (2021) [24], let \( p \in [1, +\infty) \), for any \( f \in \mathcal{H}_{\mu}^0([0, 1]^d) \) and \( \varepsilon > 0 \), there exists an FNN \( \phi \) with width

\[
N = \max\{2d \left\lceil \log_2 \left( \sqrt{d} \left( \frac{3\mu}{\varepsilon} \right)^{1/\alpha} \right) \right\rceil, 2 \left\lceil \log_2 \frac{3\mu d^\alpha}{2\varepsilon} \right\rceil + 2\}
\]

and depth \( L = 6 \) such that \( \| \phi - f \|_{L^p([0, 1]^d)} < \varepsilon \). This FNN is constructed via activation functions chosen from the set \( \{\sin(x), \max\{0, x\}, 2^x\} \). Similar to the proof for Thm. 2.4, there are \( N(d + 1) + (L - 1)N^2 \) operators “\( \times \)”, \( Nd - 1 + (L - 1)N(N - 1) \) operators “\(+\)”, and \( NL \) evaluations of activation functions to evaluate an output of the FNN. Therefore, the total number of operators in \( \phi \) as a mathematical expression is \( \mathcal{O}(d^2(\log d + \log \frac{1}{\varepsilon})^2) \), which completes the proof. \( \square \)

6.2. Algorithm to solve CO with reinforcement learning. We have included the pseudo-code for the proposed FEX implementation in Alg. 1 and Alg. 2. Specifically, Alg. 1 outlines the procedure for solving the CO problem using a predefined fixed tree. On the other hand, Alg. 2 is designed to iteratively expand a tree, thereby enhancing its expressiveness in the pursuit of identifying an improved solution based on the approach described in Alg. 1.

6.3. Using a fully connected network to model a controller. We illustrate the use of a fully connected neural network to model the controller, which generates probability mass functions for each node within a tree, in Figure 6.1. As an example, let’s assume we have a tree with three nodes: \( n_1 = 2 \) nodes for the binary set with \(|\mathbb{B}| = 2\) operators, and \( n_2 = 1 \) node for the unary set with \(|\mathbb{U}| = 3\) operators. Consequently, the output size of the neural network is \( n_1|\mathbb{B}| + n_2|\mathbb{U}| = 7 \). This output is divided into \( n_1 + n_2 \) parts, each corresponding to the probability mass function for its respective node. Besides, the input to the neural network is simply a constant vector.

6.4. Iterative method for eigenpairs. First, by solving Eqn. (4.8) with our FEX, we obtain an estimated eigenfunction \( u(x; \mathcal{T}, \mathbf{\hat{e}}, \mathbf{\hat{\theta}}) \) and get the initial estimation of the eigenvalue through the Rayleigh quotient \( \gamma_0 = \mathcal{I}(u(\cdot; \mathcal{T}, \mathbf{\hat{e}}, \mathbf{\hat{\theta}})) \). Then we can utilize Eqn. (4.9) to iteratively find the eigenpair. We define the function for Eqn. (4.9) by

\[
(6.1) \quad \mathcal{L}_2(v, \gamma) := \| - \Delta v - \| \nabla v \|_2^2 + \| x \|_2^2 - \gamma \|_{L^2(\Omega)}^2.
\]

Given \( \gamma_i \), we aim to find \( v \) that is expressed by mathematical expression and minimizes \( \mathcal{L}_2(v, \gamma_i) \). Assume \( v \) is expressed by a binary tree \( (v := v(\cdot; \mathcal{T}, \mathbf{\hat{e}}, \mathbf{\hat{\theta}})) \), and then we can search the solution
Algorithm 1 FEX with a fixed tree

**Input:** PDE and the associated functional $\mathcal{L}$; A tree $\mathcal{T}$; Searching loop iteration $T$; Coarse-tune iteration $T_1$ with Adam; Coarse-tune iteration $T_2$ with BFGS; Fine-tune iteration $T_3$ with Adam; Pool size $K$; Batch size $N$.

**Output:** The solution $u(x; \mathcal{T}, \hat{e}, \hat{\theta})$.

1: Initialize the agent $\chi$ for the tree $\mathcal{T}$
2: $\mathbb{P} \leftarrow \{\}$
3: for $T$ from 1 to $T$ do
4: Sample $N$ sequences $\{e^{(1)}, e^{(2)}, \ldots, e^{(N)}\}$ from $\chi$
5: for $n$ from 1 to $N$ do
6: Optimize $\mathcal{L}(u(x; \mathcal{T}, e^{(n)}))$ through coarse-tuning over $T_1 + T_2$ iterations to get score $S(e^{(n)})$
7: if $e^{(n)}$ belongs to the top-$K$ of all scorings in $\mathbb{P}$ then
8: $\mathbb{P}.$append($e^{(n)}$)
9: $\mathbb{P}$ pops some $e$ of the smallest score when $|\mathbb{P}| > K$
10: end if
11: end for
12: Update $\chi$ using (3.9)
13: end for
14: for $e$ in $\mathbb{P}$ do
15: Fine-tune $\mathcal{L}(u(x; \mathcal{T}, e, \theta))$ with $T_3$ iterations.
16: end for
17: return the expression with the smallest fine-tune error.

Algorithm 2 FEX with progressively expanding trees

**Input:** Tree set $\{T_1, T_2, \cdots\}$; Error tolerance $\epsilon$;

**Output:** the solution $u(x; \hat{T}, \hat{e}, \hat{\theta})$.

1: for $T$ in $\{T_1, T_2, \cdots\}$ do
2: Initialize the agent $\chi$ for the tree $T$
3: Obtain $u(x; T, \hat{e}, \hat{\theta})$ from Algorithm 1
4: if $\mathcal{L}(u(\cdot; \mathcal{T}, \hat{e}, \hat{\theta})) \leq \epsilon$ then
5: Break
6: end if
7: end for
8: return the expression with the smallest functional value.

using our FEX with the following optimization,

\[(6.2) \quad e^*_i, \theta^*_i \approx \arg \min_{e, \theta} \mathcal{L}_2(v(\cdot; \mathcal{T}, e, \theta), \gamma_i).\]

Next, we can compute the current estimated eigenvalue by $\gamma_{i+1} = \mathcal{I}(\exp(v(\cdot; \mathcal{T}, e^*_i, \theta^*_i)))$.

If continuing this loop for $G$ times, we will obtain the eigenpair $\gamma_G$ and $\exp(v(\cdot; \mathcal{T}, e^*_G, \theta^*_G))$.

**Implementation.** In our iteration method, the number of the iterative loop is $G = 10$. $\lambda_1 = \lambda_2 = 500$ in (4.8). The batch size for estimating the first term and third term of (4.8) is 10,000 while that of the second term (boundary) is 2,000. ReLU$^2$ is used in ResNet, following [14].
Fig. 4.1: Distribution of the node values and the error comparison. In FEX, we search the optimal sequence of node values and show the frequency of the node values of the binary tree in the candidate pool, consisting of the root (Root), middle node (Middle), and two leaves (Leaf 1 and Leaf 2). Based on the observation of the distribution, we readily design the new NN parameterization (FEX-NN) to estimate the solution. The last column displays comparison of the relative $L_2$ error as the function of the dimension over 6 independent trials between FEX, NN, FEX-NN and Genetic Programming (GP) for various high-dimensional PDE problems. Rows (a), (b) and (c) represent the results for Poisson equation (4.1), Linear conservation law (4.2) and Nonlinear Schrödinger equation (4.4) respectively. For various dimensions, FEX identifies the true solution, approximating solutions of almost the machine accuracy.
Fig. 4.2: Relative $L_2$ error of solutions estimated by trees with increasing depth for the problems of various dimensions. We exclude the square operator $(\cdot)^2$ in the unary set $U$, and the binary tree defined in Sec. 3.1 can not reproduce the true solution (sum of the square of coordinates) exactly in the example of the Poisson equation. We found that smaller errors could be obtained with larger tree sizes.
Find the smallest eigenvalue $\gamma$

$$\begin{align*}
-\Delta u + w \cdot u &= \gamma u, x \in \Omega \\
u &= 0, x \in \partial \Omega
\end{align*}$$

FEX

Solution

Observation

Assume

Simplification

New iterative solver

Root node

| cos | sin | exp | pow(4) | pow(3) | pow(2) | log | 1 |
|-----|-----|-----|--------|--------|--------|-----|---|
|     |     |     |        |        |        |     |   |
|     |     |     |        |        |        |     |   |
|     |     |     |        |        |        |     |   |
|     |     |     |        |        |        |     |   |

Fig. 4.3: Eigenvalue problems and postprocessing algorithm design with FEX. Bottom left: We observed that the exponent operator “$\exp(\cdot)$” dominates the tree root in the FEX searching loop. Based on this observation, we assume the solution is $\exp(v(x))$ and simplify the original PDE to a new PDE that avoids the trivial solution. Bottom right: The NN-based method produces a large error on the eigenvalue estimation, especially when the dimension is high ($d = 10$). With the postprocessing algorithm with FEX, we can identify the eigenvalue with an error close to zero.
Fig. 6.1: Illustration of using a fully connected neural network to model the controller that outputs the probability mass functions for each of the node within a tree. As an example, the tree contains three nodes: \( n_1 = 2 \) nodes for the binary set with \( |\mathbb{B}| = 2 \) operators and \( n_2 = 1 \) nodes for the unary set with \( |\mathbb{U}| = 3 \) operators.