Rating Alternatives from Pairwise Comparisons by Solving Tropical Optimization Problems

N. Krivulin

Abstract

We consider problems of rating alternatives based on their pairwise comparison under various assumptions, including conditions on the final scores of alternatives. The problems are formulated in the framework of tropical mathematics to approximate pairwise comparison matrices by reciprocal matrices of unit rank. We consider unconstrained and constrained approximation of one matrix, and simultaneous approximation of several matrices. The approximation problems are written in a common form for both multiplicative and additive comparison scales. To solve the problems, we apply recent results in tropical optimization, which provide new complete direct solutions given in a compact vector form. The new solutions extend known results and involve less computational effort. As an illustration, numerical examples of rating alternatives are presented.

Key-Words: tropical mathematics, idempotent semifield, constrained optimization, matrix approximation, reciprocal matrix, pairwise comparisons, rating of alternatives.

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1 Introduction

Tropical (idempotent) mathematics, which studies idempotent semirings [1 2 3 11 6 7 8], finds increasing application in solving real-world problems in various fields, including decision making. To solve these problems in the framework of tropical mathematics, they are formulated as optimization problems to minimize or maximize functions defined on vectors over idempotent semifields (see, e.g., an overview in [9]).

One of the applications of tropical optimization is concerned with the analysis of preferences by using pairwise comparison data in decision making. A problem of rating alternatives on the basis of their pairwise comparison

Faculty of Mathematics and Mechanics, Saint Petersburg State University, 28 Universitetsky Ave., St. Petersburg, 198504, Russia, nkk@math.spbu.ru.
matrix is examined in [10, 11, 12]. The problem is solved as an optimization problem in terms of the tropical mathematics. The vector of final scores of alternatives is found as a tropical eigenvector of the matrix.

In this paper, we offer new solutions to the problems of rating alternatives under various assumptions, including conditions on the final scores of alternatives. The problems are formulated in the framework of tropical optimization to approximate pairwise comparison matrices by reciprocal matrices of unit rank. We consider unconstrained and constrained approximation of one matrix, and simultaneous approximation of several matrices. The approximation problems are written in a common form for both multiplicative and additive comparison scales. To solve the problems, we apply recent results in [13, 14, 15], which provide new complete direct solutions given in a compact vector form. The new solutions extend known results and involve less computational effort. As an illustration, numerical examples of rating alternatives from pairwise comparison matrices are presented.

2 Rating Alternatives from Pairwise Comparison

Pairwise comparison techniques are widely used to obtain and arrange source data in the analysis of preferences in decision making (see, e.g., [16, 17, 18]). Given results of pairwise comparison of alternatives on an appropriate scale, the analysis focuses on forming judgment on the overall preference of each alternative by evaluating its individual rating (score, priority).

2.1 Pairwise Comparison Matrices

The results of comparing alternatives in pairs with multiplicative or additive scales are described by pairwise comparison matrices that have a specific antisymmetric form. Let \( A = (a_{ij}) \) be a pairwise comparison matrix. If the matrix is obtained on the basis of a multiplicative scale, then each entry \( a_{ij} \) shows that alternative \( i \) is preferred to \( j \) by \( a_{ij} \) times. The multiplicative comparison matrix is reciprocal, which means that its entries are positive and satisfy the condition

\[
a_{ij} = 1/a_{ji}.
\]

In the case of additive scale, the entry \( a_{ij} \) in the matrix \( A \) indicates by how many score units the preference of \( i \) is greater than that of \( j \). Then, the matrix \( A \) is skew-symmetric with the entries that answer the equality

\[
a_{ij} = -a_{ji}.
\]

Note that, in practice, the matrices composed of the results from pairwise comparisons on a multiplicative (additive) scale can have a form that is not reciprocal (skew-symmetric), and thus need corrections.
To provide consistency and interpretability of the preference relation, the results of pairwise comparison have to be transitive, which implies that the entries of the multiplicative (additive) comparison matrix comply with the equality

\[ a_{ij} = a_{ik}a_{kj}, \quad (a_{ij} = a_{ik} + a_{kj}). \]

A pairwise comparison matrix with transitive entries is called consistent. Every consistent matrix has a well-known form defined by a vector. Specifically, for any multiplicative (additive) consistent matrix \( A = (a_{ij}) \), there is a vector \( \mathbf{x} = (x_i) \) that completely specifies the entries of \( A \) as follows:

\[ a_{ij} = \frac{x_i}{x_j}, \quad (a_{ij} = x_i - x_j). \]

At the same time, if a matrix \( A \) is consistent, then its corresponding vector \( \mathbf{x} \) can be considered to directly represent the individual overall scores of alternatives, and thus provides the result, which is needed in the analysis of preference.

### 2.2 Approximation by Consistent Matrices

The matrices of pairwise comparison, which appear in real-world applications, are generally inconsistent, and may even have a non-antisymmetric form, which is due to various reasons, from limitations in human judgment to data errors. This leads to the problem of approximating a matrix \( A \) obtained from pairwise comparisons by consistent matrices \( \mathbf{X} \),

\[
\text{minimize } \rho(A, \mathbf{X}),
\]

where the minimum is taken over all consistent matrices \( \mathbf{X} \), and \( \rho \) is a suitable measure of approximation error.

Since the entries of any consistent matrix \( \mathbf{X} \) is uniquely determined by the elements of a vector \( \mathbf{x} \), problem (1) is equivalent to finding this vector. Considering that the vector \( \mathbf{x} \) shows the overall individual ratings of alternatives, the evaluation of preferences is reduced to solving (1).

Several approaches exist to solve problem (1), including approximation with the principal eigenvector of the matrix \( A \) [17, 19], least squares approximation [19, 20] and other techniques [21, 22, 23]. As a rule, these approaches offer algorithmic solutions using iterative numerical procedures, such as power iterations in the principal eigenvector method and the Newton algorithm in the least squares approximation.

Another approach based on tropical mathematics is proposed and examined in [10, 11, 12]. This approach uses the approximation by a consistent matrix defined by a tropical eigenvector, and hence can be considered a tropical counterpart of the conventional principal eigenvector method. Moreover, it is shown in [11] that the matrix, which solves the approximation problem, can be given not only by tropical eigenvectors, but also by some other ones.
A technique to find these vectors is proposed, which offers a computational algorithm, rather than provides a direct solution in an explicit form.

Below, we formulate the problem of finding an approximate consistent matrix as a general problem of approximation by reciprocal matrices of rank 1 in the topical mathematics sense. We show how results in tropical optimization can be applied to provide a complete direct solution to the problem, and give numerical examples.

3 Preliminary Definitions and Remarks

In this section, we outline preliminary definitions and results of tropical mathematics from [13, 9, 14, 15] to provide an appropriate analytical framework for the solutions in the subsequent sections. Further details at both introductory and advanced levels can be found, for instance, in [1, 2, 3, 4, 5, 6, 7, 8].

3.1 Idempotent Semifield

Let $\mathbb{X}$ be a set with two distinct elements $0$ and $1$, called the zero and the unit, and two binary operations $\oplus$ and $\otimes$, called addition and multiplication, such that $(\mathbb{X}, \oplus, 0)$ is an idempotent commutative monoid, $(\mathbb{X}, \otimes, 1)$ is an abelian group, multiplication distributes over addition, and the zero is absorbing for multiplication. Under these conditions, the system $(\mathbb{X}, \oplus, \otimes, 0, 1)$ is referred to as the idempotent semifield.

In the semifield, addition is idempotent to have $x \oplus x = x$ for all $x \in \mathbb{X}$. Multiplication is invertible, which means that each nonzero $x \in \mathbb{X}$ has its inverse $x^{-1}$ such that $x \otimes x^{-1} = 1$. The integer powers represent iterated products as $x^0 = 1$, $x^p = x^{p-1}x$, $x^{-p} = (x^{-1})^p$ for any $x \neq 0$ and $p > 0$.

The semifield is assumed to have a linear order that is consistent with the partial order induced by idempotent addition to define $x \leq y$ if and only if $x \oplus y = y$. Moreover, the semifield is considered algebraically closed (radicable), which means that the equation $x^p = a$ has solutions for any $a \in \mathbb{X}$ and integer $p > 0$, and thus allows the powers with rational exponents to be defined as well.

In the expressions that follow, the multiplication sign is usually omitted for the sake of brevity.

Examples of the idempotent semifield under study include

$$
\mathbb{R}_{\text{max}, \times} = (\mathbb{R}_+ \cup \{0\}, \max, \times, 0, 1),
\mathbb{R}_{\text{max}, +} = (\mathbb{R} \cup \{-\infty\}, \max, +, -\infty, 0),
$$

where $\mathbb{R}$ is the set of real numbers, and $\mathbb{R}_+ = \{x \in \mathbb{R} | x > 0\}$. The semifield $\mathbb{R}_{\text{max}, \times}$ has the addition $\oplus$ defined as maximum, and the multiplication $\otimes$ defined as usual. The neutral elements $0$ and $1$ coincide.
with the arithmetic zero and one. The power and inversion notation has the standard meaning.

The semifield \( \mathbb{R}_{\text{max,+}} \) is equipped with \( \oplus = \text{max}, \quad \otimes = +, \quad 0 = -\infty \) and \( 1 = 0 \). For each \( x \in \mathbb{R} \), the inverse \( x^{-1} \) coincides with the usual opposite number \(-x\). For all \( x, y \in \mathbb{R} \), the power \( x^y \) corresponds to the regular arithmetic product \( xy \).

In both semifields, the idempotent addition induces the order, which is consistent with the natural linear order on \( \mathbb{R} \).

### 3.2 Vector and Matrix Algebra

The set of column vectors with \( n \) elements over \( \mathbb{X} \) is denoted \( \mathbb{X}^n \). A vector with all elements equal to \( 0 \) is the zero vector. A vector is called regular if it has no zero elements.

Vector addition and scalar multiplication follow the conventional element-wise rules, where the scalar operations \( \oplus \) and \( \otimes \) play the roles of the standard addition and multiplication.

A vector \( b \) is linearly dependent on vectors \( a_1, \ldots, a_m \) if \( b = x_1 a_1 \oplus \cdots \oplus x_m a_m \) for some scalars \( x_1, \ldots, x_m \). Vectors \( a \) and \( b \) are collinear if \( b = x a \) for some scalar \( x \).

The system of vectors \( a_1, \ldots, a_m \) is linearly dependent if at least one of them is dependent on others, and independent otherwise. The set of linear combinations \( x_1 a_1 \oplus \cdots \oplus x_m a_m \) for all possible coefficients \( x_1, \ldots, x_m \) is closed under vector addition and scalar multiplication, and is referred to as the idempotent vector space generated by the system.

For each nonzero column vector \( x = (x_i) \), the multiplicative conjugate transpose is the row vector \( x^- = (x^-_i) \) with the elements \( x^-_i = x_i^{-1} \) if \( x_i \neq 0 \), and \( x^-_i = 0 \) otherwise.

The matrices with \( m \) rows and \( n \) columns form the set \( \mathbb{X}^{m \times n} \). A matrix with all entries equal to \( 0 \) is the zero matrix.

Matrix addition, matrix multiplication and scalar multiplication are routinely defined entry-wise, where the operations \( \oplus \) and \( \otimes \) are used instead of the usual addition and multiplication.

For any nonzero matrix \( A = (a_{ij}) \in \mathbb{X}^{m \times n} \), the multiplicative conjugate transpose is the matrix \( A^- = (a^-_{ij}) \in \mathbb{X}^{n \times m} \) with the entries \( a^-_{ij} = a^{-1}_{ji} \) if \( a_{ji} \neq 0 \), and \( a^-_{ij} = 0 \) otherwise.

The rank of a matrix is defined as the maximum number of linearly independent columns (rows) in the matrix. A matrix \( A \) has rank 1 if and only if there exist nonzero column vectors \( x \) and \( y \) such that \( A = xy^T \).

Consider square matrices of order \( n \) in the set \( \mathbb{X}^{n \times n} \). A matrix that has 1 along the diagonal, and 0 elsewhere, is the identity matrix denoted \( I \). The power notation is defined to indicate repeated multiplication as \( A^0 = I \) and \( A^p = A^{p-1} A \) for any square matrix \( A \) and integer \( p > 0 \).
A matrix $A$ without zero entries is called symmetrically reciprocal (or, simply, reciprocal) if the condition $A^T = A$ holds. A reciprocal matrix $A$ is of unit rank if and only if $A = xx^T$, where $x$ is a regular column vector.

The trace of a matrix $A = (a_{ij})$ is given by
\[
\text{tr} A = a_{11} \oplus \cdots \oplus a_{nn}.
\]

For any matrices $A$ and $B$, the following equalities hold:
\[
\text{tr}(A \oplus B) = \text{tr} A \oplus \text{tr} B, \quad \text{tr}(AB) = \text{tr}(BA).
\]

### 3.3 Distance Functions

The distance between two regular vectors $x, y \in \mathbb{X}^n$ is given by the function
\[
\rho(x, y) = y^T x \oplus x^T y,
\]
which takes the minimum value $1$ only when $y = x$.

For the real semifield $\mathbb{R}_{\text{max,}+}$, where $1 = 0$, the function $\rho$ coincides with the Chebyshev metric
\[
\rho_\infty(x, y) = \max_{1 \leq i \leq n} |x_i - y_i|.
\]

In the case of $\mathbb{R}_{\text{max,}x}$, the function $\rho$ differs from the usual metrics in the range of values, and becomes a log-Chebyshev metric after taking the logarithm. In the general case, the function $\rho$ is referred to as the Chebyshev-like distance.

The distance between two matrices without zero entries is defined by the Chebyshev-like distance function
\[
\rho(A, B) = \text{tr}(B^T A) \oplus \text{tr}(A^T B).
\]

This function has the form of the Chebyshev metric for the semifield $\mathbb{R}_{\text{max,}+}$, and takes the form of a log-Chebyshev metric after logarithmic transformation for $\mathbb{R}_{\text{max,}x}$.

### 3.4 Eigenvalues and Eigenvectors of Matrices

A scalar $\lambda \in \mathbb{X}$ is an eigenvalue of a matrix $A \in \mathbb{X}^{n \times n}$ if there exists a nonzero vector $x \in \mathbb{X}^n$ such that $Ax = \lambda x$. This vector $x$ is an eigenvector of $A$, corresponding to $\lambda$.

The maximum eigenvalue of a matrix $A = (a_{ij})$ is referred to as the spectral radius of the matrix and calculated as
\[
\lambda = \text{tr} A \oplus \cdots \oplus \text{tr}^{1/n}(A^n),
\]
or, in terms of matrix entries, as
\[
\lambda = \bigoplus_{k=1}^n \bigoplus_{1 \leq i_1, \ldots, i_k \leq n} (a_{i_1i_2}a_{i_2i_3} \cdots a_{i_ki_1})^{1/k}.
\]
4 Tropical Optimization Problems

We now consider optimization problems that are formulated and solved in the tropical mathematics setting. To represent complete direct solutions to the problems, we introduce a function that takes any matrix $A \in X^{n \times n}$ to produce the scalar

$$\text{Tr}(A) = \text{tr} A \oplus \cdots \oplus \text{tr} A^n.$$ 

Provided that $\text{Tr}(A) \leq 1$, we calculate the matrix (also known as the Kleene star)

$$A^* = I \oplus A \oplus \cdots \oplus A^{n-1}.$$ 

We start with the unconstrained optimization problem: given a matrix $A \in X^{n \times n}$, find regular vectors $x \in X^n$ that minimize

$$x - Ax,$$ 

(4)

A complete direct solution to the problem is provided in [13, 14, 15] in the following form.

**Lemma 1.** Let $A$ be a matrix with spectral radius $\lambda > 0$, and $A_\lambda = \lambda^{-1} A$. Then, the minimum value in problem (4) is equal to $\lambda$, and all regular solutions are given by

$$x = A_\lambda^* u, \quad u \in X^n.$$ 

It follows from Lemma 1 that the solutions form an idempotent vector space spanned by the columns of the matrix $A_\lambda^*$. Suppose now that, given matrices $A, C \in X^{n \times n}$, we need to find regular solutions $x \in X^n$ to the problem

$$\text{minimize } x - Ax,$$ 

subject to $Cx \leq x$. 

(5)

The next complete solution to the problem is given in [14].

**Theorem 2.** Let $A$ be a matrix with spectral radius $\lambda > 0$, and $C$ a matrix with $\text{Tr}(C) \leq 1$. Then, the minimum value in problem (5) is equal to

$$\theta = \lambda \oplus \bigoplus_{k=1}^{n-1} \bigoplus_{1 \leq i_1 + \cdots + i_k \leq n-k} \text{tr}^{1/k}(AC_{i_1} \cdots AC_{i_k}),$$

and all regular solutions are given by

$$x = (\theta^{-1} A \oplus C)^* u, \quad u \in X^n.$$ 

In the ensuing section, the above solutions are applied to solve matrix approximation problems, which appear in rating alternatives on the basis of their pairwise comparisons.
5 Evaluation of Scores by Pairwise Comparisons

We are now in a position to put the problem of rating alternatives via approximation by consistent matrices in the context of tropical optimization. First, we note that, in the framework of tropical mathematics, both multiplicative and additive consistent matrices can be represented in a common form of the reciprocal matrix of rank 1, which are given by

\[ X = xx^-, \]

to be interpreted in terms of either the semifield \( \mathbb{R}_{\text{max}, \times} \) for the multiplicative case and the semifield \( \mathbb{R}_{\text{max}, +} \) for the additive.

The problem of finding an approximate consistent matrix \( X \), or, equivalently, a vector of scores \( x \), can be described as follows. Given a matrix \( A \in \mathbb{R}^{n \times n} \), find regular vectors \( x \in \mathbb{R}^n \) that minimize

\[ \rho(A, xx^-), \]

where \( \rho \) is a measure of approximation error, which is given by the Chebyshev-like distance function defined as [2]. The function \( \rho \) becomes a log-Chebyshev metric for multiplicative scale, and the Chebyshev metric for the additive.

In this section, we apply the solutions of tropical optimization problems given by Lemma 1 and Theorem 5 to directly solve approximation problem (6). The results obtained are then used to evaluate, under various assumptions, the scores of alternatives, based on pairwise comparison matrices.

These results offer complete direct solutions given in a compact vector form, which extend the known solutions to the problems in [10, 11], and are easier to calculate.

Below, we examine problems of rating alternatives on a multiplicative scale. Considering that, in terms of tropical mathematics, the solution to approximation problems in multiplicative and additive settings has a common general form, the case of additive scale is not covered here for brevity.

5.1 Evaluation of Scores Given by One Matrix

We first provide a solution for evaluating the vector of scores \( x \) on the basis of one matrix \( A \), which represents the results of pairwise comparisons. The problem is formulated in the setting of the semifield \( \mathbb{R}_{\text{max}, \times} \) in the form of (6) to approximate the matrix \( A \) by a reciprocal matrix of unit rank.

**Theorem 3.** Let \( A \) be a matrix such that the matrix \( B = A \oplus A^- \) has no zero entries, \( \mu \) be the spectral radius of \( B \), and \( B_\mu = \mu^{-1}B \). Then the minimum value in problem (6) is equal to \( \mu \) and all solutions are given by

\[ x = B_\mu^*u, \quad u \in \mathbb{R}^n_+. \]
Proof. It is easy to see that, since all entries in the matrix $B$ are nonzero, this matrix has the spectral radius $\mu > 0$.

We write the objective function by using (2). It follows from the equality $(xx^-) = xx^-$ and properties of the trace that

$$\rho(A, xx^-) = \text{tr}((xx^-) A) + \text{tr}(A^- xx^-) = x^- A x + x^- A^- x = x^- B x.$$ 

An application of Lemma 1 completes the proof. \qed

We now give an example of evaluating the score vector from a reciprocal matrix of pairwise comparisons. For arbitrary positive matrices, evaluation of scores follows the same way.

Example 1. Consider the reciprocal matrix defined as

$$A = \begin{pmatrix} 1 & 3 & 2 & 4 \\ 1/3 & 1 & 1/3 & 1/2 \\ 1/2 & 3 & 1 & 1/4 \\ 1/4 & 2 & 4 & 1 \end{pmatrix}.$$ 

To approximate the matrix by a reciprocal matrix of unit rank, and thus to find a score vector $x$, we apply Theorem 3.

Since the matrix $A$ is reciprocal, and hence $A^{-} = A$, we see that $B = A \oplus A^{-} = A$. We apply (3) to find the spectral radius of the matrix $B$ to be $\mu = 2$. Furthermore, we take the matrix

$$B_\mu = \mu^{-1}B = \begin{pmatrix} 1/2 & 3/2 & 1 & 2 \\ 1/6 & 1/2 & 1/6 & 1/4 \\ 1/4 & 3/2 & 1/2 & 1/8 \\ 1/8 & 1 & 2 & 1/2 \end{pmatrix},$$

and then calculate the powers

$$B_\mu^2 = \begin{pmatrix} 1/4 & 2 & 4 & 1 \\ 1/12 & 1/4 & 1/2 & 1/3 \\ 1/4 & 3/4 & 1/4 & 1/2 \\ 1/2 & 3 & 1 & 1/4 \end{pmatrix}, \quad B_\mu^3 = \begin{pmatrix} 1 & 6 & 2 & 1/2 \\ 1/8 & 3/4 & 2/3 & 1/6 \\ 1/8 & 1/2 & 1 & 1/2 \\ 1/2 & 3/2 & 1/2 & 1 \end{pmatrix}.$$ 

Finally, we compose the matrix

$$B_\mu^* = I \oplus B_\mu \oplus B_\mu^2 \oplus B_\mu^3 = \begin{pmatrix} 1 & 6 & 4 & 2 \\ 1/6 & 1 & 2/3 & 1/3 \\ 1/4 & 3/2 & 1 & 1/2 \\ 1/2 & 3 & 2 & 1 \end{pmatrix}.$$ 

Clearly, the columns in the matrix $B_\mu^*$ are collinear to each other. Specifically, the last three columns can be obtained by multiplying the first one
by 6, 4 and 2, respectively. Since each column generates exactly the same vector space, it is sufficient to take only one of them to describe all solution vectors \( \mathbf{x} \). We use the first column and write the score vector as

\[
\mathbf{x} = \begin{pmatrix} 1 \\ 1/6 \\ 1/4 \\ 1/2 \end{pmatrix} u,
\]

where \( u \) is an arbitrary positive number to be fixed in accordance with the required form or interpretation of the result.

With \( u = 1 \), the vector \( \mathbf{x} = (1, 1/6, 1/4, 1/2)^T \) shows that the first alternative is of the highest score \( \mathbf{x}_1 = 1 \), followed by the fourth and third with scores \( \mathbf{x}_4 = 1/2 \) and \( \mathbf{x}_3 = 1/4 \). The second alternative has the lowest score \( \mathbf{x}_2 = 1/6 \).

If the scores are considered as weights, which must add up to one, we put \( u = 1/(1 + 1/6 + 1/4 + 1/2) = 12/23 \). The vector takes the form \( \mathbf{x} = (12/23, 2/23, 3/23, 6/23)^T \).

### 5.2 Evaluation of Scores Given by Several Matrices

Suppose that there are \( m \) matrices \( A_1, \ldots, A_m \in \mathbb{K}^{n \times n} \), and we need to determine a reciprocal matrix of rank 1 that approximates these matrices simultaneously. The approximation problem is formulated in terms of the semifield \( \mathbb{K}_{\text{max}} \times \mathbb{K}_{\text{max}} \) in the form of (6) to find regular vectors \( \mathbf{x} \) that

\[
\min \left( \max_{1 \leq i \leq m} \rho(\mathbf{A}_i, \mathbf{x} \mathbf{x}^{-}) \right).
\]

#### Theorem 4

Let \( \mathbf{A}_i \) be matrices for all \( i = 1, \ldots, m \) such that the matrix \( \mathbf{B} = \mathbf{A}_1 \oplus \mathbf{A}_1^{-} \oplus \cdots \oplus \mathbf{A}_m \oplus \mathbf{A}_m^{-} \) has no zero entries, \( \mu \) be the spectral radius of \( \mathbf{B} \), and \( \mathbf{B}_u = \mu^{-1} \mathbf{B} \). Then, the minimum value in problem (7) is equal to \( \mu \) and all solutions are given by

\[
\mathbf{x} = \mathbf{B}_u^* \mathbf{u}, \quad \mathbf{u} \in \mathbb{R}_+^n.
\]

**Proof.** For each \( i = 1, \ldots, m \), we use the same argument as in Theorem 3 to write \( \rho(\mathbf{A}_i, \mathbf{x} \mathbf{x}^{-}) = \mathbf{x}^{-} (\mathbf{A}_i \oplus \mathbf{A}_i^{-}) \mathbf{x} \). Then, we represent the objective function as

\[
\max_{1 \leq i \leq m} \rho(\mathbf{A}_i, \mathbf{x} \mathbf{x}^{-}) = \bigoplus_{i=1}^{m} \mathbf{x}^{-} (\mathbf{A}_i \oplus \mathbf{A}_i^{-}) \mathbf{x} = \mathbf{x}^{-} \mathbf{B} \mathbf{x}.
\]

The desired result now follows from Lemma 1. \( \square \)
Example 2. We now evaluate the score vector based on the simultaneous approximation of \( m = 2 \) reciprocal matrices

\[
A_1 = \begin{pmatrix}
1 & 3 & 2 & 4 \\
1/3 & 1 & 1/3 & 1/2 \\
1/2 & 3 & 1 & 1/3 \\
1/4 & 2 & 3 & 1
\end{pmatrix}, \quad A_2 = \begin{pmatrix}
1 & 4 & 2 & 3 \\
1/4 & 1 & 1/2 & 1/2 \\
1/2 & 2 & 1 & 1/4 \\
1/3 & 2 & 4 & 1
\end{pmatrix}.
\]

To solve the problem by applying Theorem 4, we have to compose the matrix 
\( B = A_1 \oplus A_1^{-1} \oplus A_2 \oplus A_2^{-1} \). Considering that the matrices \( A_1 \) and \( A_2 \) are reciprocal, we have

\[
B = A_1 \oplus A_2 = \begin{pmatrix}
1 & 4 & 2 & 4 \\
1/3 & 1 & 1/2 & 1/2 \\
1/2 & 3 & 1 & 1/3 \\
1/3 & 2 & 4 & 1
\end{pmatrix}.
\]

Note that the obtained matrix coincides with the matrix \( B \) in Example 1. Since the matrix \( B \) completely determines the set of solution vectors, we can use the result of this example, which offers the score vector \( x = (1, 1/6, 1/4, 1/2)^T \).

5.3 Constrained Evaluation of Scores

Consider the problem of evaluating the vector \( x = (x_i) \), which represents the individual overall scores calculated from the results of pairwise comparison given by a matrix \( A \). Suppose that, for some reasons, additional constraints are imposed on the scores by inequalities in the form \( x_i \geq c_{ij}x_j \), which requires that the overall score of alternative \( i \) must be \( c_{ij} \) times greater or more than the score of alternative \( j \).

To describe the problem in terms of tropical mathematics, we introduce a matrix \( C = (c_{ij}) \), where we put \( c_{ij} = 0 \) if no constraint is defined for alternatives \( i \) and \( j \). It is easy to see that the constraints can be represented as the inequality \( Cx \leq x \) written in terms of the semifield \( \mathbb{R}_{\text{max},x} \).

By combining the constraint with the objective function in the framework of \( \mathbb{R}_{\text{max},x} \), we arrive at the next constrained approximation problem. Given matrices \( A, C \in \mathbb{R}^{n \times n} \), the problem is to find regular vectors \( x \) that

\[
\begin{align*}
\text{minimize} & \quad \rho(A, xx^-) \\
\text{subject to} & \quad Cx \leq x.
\end{align*}
\]

\[ (8) \]

Theorem 5. Let \( A \) be a matrix such that the matrix \( B = A \oplus A^- \) has no zero entries, \( \mu \) be the spectral radius of \( B \), and \( C \) a matrix with \( \text{Tr}(C) \leq 1 \). Then, the minimum value in problem (3) is equal to

\[
\theta = \mu \oplus \bigoplus_{k=1}^{n-1} \bigoplus_{1 \leq i_1 + \cdots + i_k \leq n-k} \text{tr}^{1/k}(BC_{i_1} \cdots BC_{i_k}),
\]
and all regular solutions are given by
\[ x = (\theta^{-1}B \oplus C)^* u, \quad u \in \mathbb{R}^n. \]

**Proof.** To verify the statement, we rewrite the objective function as in Theorem 3 and then apply Theorem 5.

**Example 3.** Let us evaluate scores in a constrained problem, where the results of pairwise comparison and the constraints are defined by the matrices
\[
A = \begin{pmatrix}
1 & 3 & 2 & 4 \\
1/3 & 1 & 1/3 & 1/2 \\
1/2 & 3 & 1 & 1/4 \\
1/4 & 2 & 4 & 1
\end{pmatrix}, \quad
C = \begin{pmatrix}
0 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 \\
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0
\end{pmatrix}.
\]

Note that the solution to the unconstrained problem with the matrix \( B = A \) is provided by Example 1. Furthermore, the constraints given by the matrix \( C \) take the form
\[ x_4 \leq x_2, \quad x_2 \leq x_3, \quad x_3 \leq x_4, \]
which is obviously equivalent to one condition \( x_2 = x_3 = x_4 \).

By Theorem 5 we have to calculate the value of \( \theta \). Using properties of the trace yields the expression
\[
\theta = \mu \oplus \text{tr}(BC(I \oplus C \oplus C^2)) \oplus \text{tr}^{1/2}(BCB(I \oplus C)) \oplus \text{tr}^{1/3}(BCB^2),
\]
where \( \mu = 2 \) is the spectral radius of the matrix \( B \).

We now calculate the matrices
\[
I \oplus C = \begin{pmatrix}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 1 \\
0 & 1 & 1 & 0 \\
0 & 0 & 1 & 1
\end{pmatrix}, \quad
C^2 = \begin{pmatrix}
0 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 0 & 0 & 1 \\
0 & 1 & 0 & 0
\end{pmatrix},
\]
\[
I \oplus C \oplus C^2 = \begin{pmatrix}
1 & 0 & 0 & 0 \\
0 & 1 & 1 & 1 \\
0 & 1 & 1 & 1 \\
0 & 1 & 1 & 1
\end{pmatrix}.
\]

Next, we obtain the matrices
\[
BC = \begin{pmatrix}
0 & 2 & 4 & 3 \\
0 & 1/3 & 1/2 & 1 \\
0 & 1 & 1/4 & 3 \\
0 & 4 & 1 & 2
\end{pmatrix}, \quad
BC(I \oplus C \oplus C^2) = \begin{pmatrix}
0 & 4 & 4 & 4 \\
0 & 1 & 1 & 1 \\
0 & 3 & 3 & 3 \\
0 & 4 & 4 & 4
\end{pmatrix},
\]
and then find \( \text{tr}(BC(I \oplus C \oplus C^2)) = 4 \).

12
Furthermore, we calculate

\[ BCB = \begin{pmatrix} 2 & 12 & 12 & 3 \\ 1/4 & 2 & 4 & 1 \\ 3/4 & 6 & 12 & 3 \\ 4/3 & 4 & 8 & 2 \end{pmatrix}, \quad BCB(I \oplus C) = \begin{pmatrix} 2 & 12 & 12 & 12 \\ 1/4 & 4 & 4 & 2 \\ 3/4 & 12 & 12 & 6 \\ 4/3 & 8 & 8 & 4 \end{pmatrix}, \]

from which it follows that \( \text{tr}^{1/2}(BCB(I \oplus C)) = \sqrt{12} \).

After calculating the matrix

\[ BCB^2 = \begin{pmatrix} 6 & 36 & 12 & 8 \\ 2 & 12 & 4 & 1 \\ 6 & 36 & 12 & 3 \\ 4 & 24 & 8 & 16/3 \end{pmatrix}, \]

and the trace \( \text{tr}^{1/3}(BCB^2) = \sqrt[3]{12} \), we conclude that \( \theta = 4 \).

We now calculate the matrices

\[ \theta^{-1}B \oplus C = \begin{pmatrix} 1/4 & 3/4 & 1/2 & 1 \\ 1/12 & 1/4 & 1/12 & 1 \\ 1/8 & 1 & 1/4 & 1/16 \\ 1/16 & 1/2 & 1 & 1/4 \end{pmatrix}, \]

\[ (\theta^{-1}B \oplus C)^2 = \begin{pmatrix} 1/16 & 1/2 & 1 & 3/4 \\ 1/16 & 1/2 & 1 & 1/4 \\ 1/12 & 1/4 & 1/12 & 1 \\ 1/8 & 1 & 1/4 & 1/2 \end{pmatrix}, \]

\[ (\theta^{-1}B \oplus C)^3 = \begin{pmatrix} 1/8 & 1 & 3/4 & 1/2 \\ 1/8 & 1 & 1/4 & 1/2 \\ 1/16 & 1/2 & 1 & 1/4 \\ 1/12 & 1/4 & 1/2 & 1 \end{pmatrix}. \]

Finally, consider the matrix

\[ (\theta^{-1}B \oplus C)^* = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1/8 & 1 & 1 & 1 \\ 1/8 & 1 & 1 & 1 \\ 1/8 & 1 & 1 & 1 \end{pmatrix}. \]

The last three columns assign the same score equal to one to all alternatives, and therefore, are of no interest. The first column offers a score vector \( x = (1, 1/8, 1/8, 1/8)^T \), which are consistent with both the results of pairwise comparisons, offered by Example 1 and the constraint \( x_2 = x_3 = x_4 \).

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