Exclusion and Guard Zones in DS-CDMA Ad Hoc Networks

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Abstract—The central issue in direct-sequence code-division multiple-access (DS-CDMA) ad hoc networks is the prevention of a near-far problem. This paper considers two types of guard zones that may be used to control the near-far problem: a fundamental exclusion zone and an additional CSMA guard zone that may be established by the carrier-sense multiple-access (CSMA) protocol. In the exclusion zone, no mobiles are physically present, modeling the minimum physical separation among mobiles that is always present in actual networks. Potentially interfering mobiles beyond a transmitting mobile’s exclusion zone, but within its CSMA guard zone, are deactivated by the protocol. This paper provides an analysis of DS-CSMA networks with either or both types of guard zones. A network of finite extent with a finite number of mobiles and uniform clustering as the spatial distribution is modeled. The analysis applies a closed-form expression for the outage probability in the presence of Nakagami fading, conditioned on the network geometry. The tradeoffs between exclusion zones and CSMA guard zones are explored for DS-CDMA and unspread networks. The spreading factor and the guard-zone radius provide design flexibility in achieving specified levels of average outage probability and transmission capacity. The advantage of an exclusion zone over a CSMA guard zone is that since the network is not thinned, the number of active mobiles remains constant, and higher transmission capacities can be achieved.

Index Terms—Spread-spectrum communication, Carrier-sense multiaccess, Mobile communications, Fading channels.

I. INTRODUCTION

Direct-sequence code-division multiple-access (DS-CDMA) ad hoc networks are realized by using direct-sequence spread-spectrum modulation while the mobiles or nodes of multiple users simultaneously transmit signals in the same frequency band. All signals use the entire allocated spectrum, but the spreading sequences differ. DS-CDMA is advantageous for ad hoc networks because it eliminates the need for any frequency or time-slot coordination, imposes no sharp upper bound on the number of mobiles, directly benefits from inactive terminals in the network, and is capable of efficiently implementing bursty data traffic, intermittent voice signals, multibeamed arrays, and reassignments to accommodate variable data rates. Furthermore, DS-CDMA systems are inherently resistant to interference, interception, and frequency-selective fading.

The central issue in DS-CDMA ad hoc networks is the prevention of a near-far problem. The solution to the near-far problem in cellular networks is power control. However, the absence of a centralized control of an ad hoc network renders any attempted power control local rather than pervasive. As an alternative to power control, ad hoc networks typically use guard zones surrounding mobiles to manage interference. This paper considers two types of guard zones to control the near-far problem: a fundamental exclusion zone and an additional guard zone that may be established by the carrier-sense multiple-access (CSMA) protocol; i.e., a CSMA guard zone. In the exclusion zone surrounding a mobile, no other mobiles are physically present. The exclusion zone is used to model networks for which there is a minimum physical separation among mobiles, for instance, when there is a one-to-one correspondence between mobiles and vehicles, and a spacing is required to prevent collision. In the CSMA guard zone, which extends beyond the exclusion zone, other mobiles may be present, but they are deactivated by the protocol. The CSMA guard zone offers additional near-far protection beyond that offered by the exclusion zone, but at the cost of reduced network transmission capacity, as shown subsequently. However, CSMA guard zones are useful for operating environments that do not permit large exclusion zones, such as networks with low mobility or a high density of mobile terminals.

In this paper, the network topology is modeled stochastically by specifying the spatial distribution of its mobiles. A similar approach was taken by Hasan and Andrews to study guard zones, which like leveraged tools from stochastic geometry. However, the existing literature makes the unrealistic assumption that the network is infinite in extent with mobiles that are drawn from a Poisson point process. A key weakness with the use of a Poisson point process is that it does not preclude an infinite number of mobiles from being located within a small region, which renders it inapplicable to the study of guard zones. This weakness limited to only place a guard zone around the reference receiver. In contrast, the network model in the present paper allows a guard zone around every mobile in the network, which is possible by placing minimum-spacing constraints on the spatial model. Spatial constraints are studied in in the context of infinite networks, but the analysis is intractable except when operating at high signal-to-interference ratio (SIR).

The analysis in this paper builds upon a recent closed-form expression for the exact outage probability at a reference receiver conditioned on the network realization. The analysis allows for the arbitrary placement of mobiles, arbitrary shadowing factors, fading with a Nakagami-m factor that can vary among the channels from the mobiles to the reference-receiver, thermal noise, and an arbitrary duty factor.
for each mobile. Simulation results in [14] confirmed the exactness of the analysis. In the present paper, we apply the analysis of [14] to obtain the spatially averaged outage probability for networks that use guard zones. The spatial averaging is obtained by taking the expected value of the outage probability with respect to the spatial distribution and distribution of the shadowing. While this expectation is not generally analytically tractable when guard zones are present, a simple and efficient Monte Carlo method can be used that involves the random placement of mobiles and generation of shadowing factors.

The contributions of this paper are as follows. A fundamental contribution of this paper is that it presents an effective approach to analyzing finite networks with additional spatial constraints imposed. Such constraints could include a minimum physical separation between mobiles, and the deactivation of mobiles within a certain radius of an active mobile. Using this analytical approach, the impact of both types of guard zones on performance is studied. The performance is quantified by the spatially averaged outage probability, latency, and transmission capacity. The numerical results in this paper show the influence and relationship of essential parameters including the reference receiver location, the spreading factor, the CSMA guard-zone radius, the exclusion-zone radius, the transmission distance, and the number of potentially interfering mobiles. The principal advantages of our analysis relative to ones based on stochastic geometry are the accommodation of a network of arbitrary size and shape and an arbitrary spatial distribution of mobiles, the ability to distinguish among the performances of mobiles located anywhere in the interior of a network and its outer edges, and the accommodation of arbitrary fading and shadowing parameters for every path between two mobiles.

The remainder of the paper is organized as follows. Section II describes exclusion zones and CSMA guard zones, and Section III presents the network model. Section IV reviews the conditional outage probability, where the conditioning is on the network realization. Examples of conditional outage probability both with and without CSMA guard zones are given in Section V. The conditioning is removed in Section VI by averaging the outage probability with respect to the shadowing and the spatial distribution of the mobiles. Numerical results are given in Section VII. Finally, Section VIII concludes the paper.

II. GUARD ZONES

The IEEE 802.11 standard, which is currently the predominant standard for ad hoc networks, uses CSMA with collision avoidance in its medium-access control protocol [2], [3]. The implementation entails the exchange of request-to-send (RTS) and clear-to-send (CTS) handshake packets between a transmitter and receiver during their initial phase of communication that precedes the subsequent data and acknowledgment packets. The receipt of the RTS/CTS packets with sufficient power levels by nearby mobiles causes them to inhibit their own transmissions, which would produce interference in the receiver of interest. The transmission of separate CTS packets in addition to the RTS packets decreases the possibility of subsequent signal collisions at the receiver due to nearby hidden terminals that do not sense the RTS packets. Thus, the RTS/CTS packets essentially establish guard zones surrounding a transmitter and receiver, and hence prevent a near-far problem except during the initial reception of an RTS packet. The interference at the receiver is restricted to concurrent transmissions generated by mobiles outside the guard zones. The fundamental disadvantage with this CSMA guard zone is that it inhibits other concurrent transmissions within the zone, thereby reducing the network transmission capacity.

A more fundamental guard zone, which is called the exclusion zone, is based on the spacing that occurs in actual mobile networks. For instance, when the radios are mounted on separate vehicles, there is a need for crash avoidance by maintaining a minimum vehicle separation. A small exclusion zone maintained by visual sightings exists in practical networks, but a more reliable and extensive one can be established by equipping each mobile with a global positioning system (GPS) and periodically broadcasting each mobile’s GPS coordinates. Mobiles that receive these messages could compare their locations to those in the messages and alter their movements accordingly. The major advantage of the exclusion zone compared with a CSMA guard zone is that the exclusion zone prevents near-far problems at receivers while not inhibiting any potential concurrent transmissions. Another advantage of an exclusion zone is enhanced network connectivity because of the inherent constraint on the clustering of mobiles.

When CSMA is used in a network, the CSMA guard zone will usually encompass the exclusion zone. Although both zones may cover arbitrary regions, they are modeled as circular regions in the subsequent examples for computational convenience, and the region of the CSMA guard zone that lies outside the exclusion zone is an annular ring. The existence of an annular ring enhances the near-far protection at the cost of inhibiting potential concurrent transmissions within the annular ring. The tradeoffs entailed in having the CSMA guard zone are examined subsequently.

III. NETWORK MODEL

The network comprises a fixed number of mobiles in a circular area with radius $r_{\text{net}}$, although any arbitrary two- or three-dimensional regions could be considered. A reference receiver is located inside the circle, a reference transmitter $X_0$ is located at distance $||X_0||$ from the reference receiver, and there are $M$ potentially interfering mobiles $X_1, \ldots, X_M$. The variable $X_i$ represents both the $i^{th}$ mobile and its location, and $||X_i||$ is the distance from the $i^{th}$ mobile to the reference receiver. Each mobile uses a single omnidirectional antenna. The radii of the exclusion zone and the CSMA guard zone are $r_{\text{ex}}$ and $r_{\text{g}}$, respectively.

The potentially interfering mobiles are uniformly distributed throughout the network area outside the exclusion zones, according to a uniform clustering model. One by one, the location of each $X_i$ is drawn according to a uniform distribution within the radius-$r_{\text{net}}$ circle. However, if an $X_i$ falls within the exclusion zone of a previously placed mobile, then it has a new
random location assigned to it as many times as necessary until it falls outside all exclusion zones. Unlike Matern thinning \[15\], which silences mobiles without replacing them, uniform clustering maintains a fixed number of potentially interfering mobiles. Setting the exclusion zone to \(r_{ex} = 0\) is equivalent to drawing the mobiles from a binomial point process.

Since there is no significant advantage to using short spreading sequences in an asynchronous ad hoc network, long spreading sequences are assumed and modeled as random binary sequences with chip duration \(T_c\). The spreading factor or processing gain \(G\) directly reduces the interference power. The multiple-access interference is assumed to be asynchronous, and the power from each interfering \(X_i\) is further reduced by the chip function \(h(\tau_i)\), which is a function of the chip waveform and the timing offset \(\tau_i\) of \(X_i\)’s spreading sequence relative to that of the desired signal \[1\]. Since only timing offsets modulo-\(T_c\) are relevant, \(0 \leq \tau_i < T_c\). In a network of quadruphase direct-sequence systems, a multiple-access interference signal with power \(I_i\) before despreading is reduced after despreading to the power level \(I_i h(\tau_i)/G_i\), where \[1\] and \[16\]

\[
h(\tau_i) = \frac{1}{T_c^2} [R^2(\tau_i) + R^2(T_c - \tau_i)]
\]

and \(R^2(\tau_i)\) is the partial autocorrelation for the normalized chip waveform. Thus, the interference power is effectively reduced by the factor \(G_i = G/h(\tau_i)\). Assuming a rectangular chip waveform and that \(\tau_i\) has a uniform distribution over \([0, T_c]\), the expected value of \(h(\tau_i)\) is \(2/3\).

After the despreading, the power of \(X_i\)’s signal at the reference receiver is

\[
\rho_i = \tilde{P}_i g_i 10^{\xi_i/10} f(||X_i||)
\]

(2)

where \(\tilde{P}_i\) is the received power at a reference distance \(d_0\) (assumed to be sufficiently far that the signals are in the far field) after despreading when fading and shadowing are absent, \(g_i\) is the power gain due to fading, \(\xi_i\) is a shadowing factor, and \(f(\cdot)\) is a path-loss function. The path-loss function is expressed as the power law

\[
f(d) = \left(\frac{d}{d_0}\right)^{-\alpha}, \quad d \geq d_0
\]

(3)

where \(\alpha \geq 2\) is the path-loss exponent. It is assumed that \(r_{ex} \geq d_0\). The \(\{g_i\}\) are independent with unit-mean, but are not necessarily identically distributed; i.e., the channels from the different \(\{X_i\}\) to the reference receiver may undergo fading with different distributions. For analytical tractability and close agreement with measured fading statistics, Nakagami fading is assumed, and \(g_i = a_i^2\), where \(a_i\) is Nakagami with parameter \(m_i\). When the channel between \(X_i\) and the reference receiver undergoes Rayleigh fading, \(m_i = 1\) and the corresponding \(g_i\) is exponentially distributed. In the presence of log-normal shadowing, the \(\{\xi_i\}\) are independent zero-mean Gaussian with variance \(\sigma^2_x\). For ease of exposition, it is assumed that the shadowing variance is the same for the entire network, but the results may be easily generalized to allow for different shadowing variances over parts of the network. In the absence of shadowing, \(\xi_i = 0\).

It is assumed that the \(\{g_i\}\) remain fixed for the duration of a time interval but vary independently from interval to interval (block fading). With probability \(p_i\), the \(i^{th}\) mobile transmits in the same time interval as the desired signal. The \(\{p_i\}\) can be used to model voice-activity factors, controlled silence, or failed link transmissions and the resulting retransmission attempts. The \(\{p_i\}\) need not be the same; for instance, CSMA protocols can be modeled by setting \(p_i = 0\) when a mobile lies within the CSMA guard zone of another active mobile, which is equivalent to implementing Matern thinning within the annular ring corresponding to that guard zone.

The instantaneous signal-to-interference-and-noise ratio (SINR) at the reference receiver is given by

\[
\gamma = \frac{\rho_0 M}{N + \sum_{i=1}^{M} I_i \rho_i}
\]

(4)

where \(\rho_0\) is the received power of the desired signal, \(N\) is the noise power, and the indicator \(I_i\) is a Bernoulli random variable with probability \(P[I_i = 1] = p_i\) and \(P[I_i = 0] = 1 - p_i\).

Since the despreading does not significantly affect the desired-signal power, the substitution of \[2\] and \[3\] into \[4\] yields

\[
\gamma = \frac{g_0 \Omega_0}{\Gamma^{-1} + \sum_{i=1}^{M} I_i g_i \Omega_i}
\]

(5)

where

\[
\Omega_i = \left\{ \begin{array}{ll}
10^{\xi_i/10} & \text{if } i = 0 \\
\frac{P_i}{G_i P_0} 10^{\xi_i/10} & \text{if } i > 0
\end{array} \right.
\]

(6)

is the normalized power of \(X_i\), \(P_i\) is the received power at the reference distance \(d_0\) before despreading when fading and shadowing are absent, and \(\Gamma = d_0^2 P_0/N\) is the SNR when the reference transmitter is at unit distance from the reference receiver and fading and shadowing are absent.

IV. OUTAGE PROBABILITY

The outage probability quantifies the likelihood that the noise and interference will be too severe for useful communications. Outage probability is defined with respect to an SINR threshold \(\beta\), which represents the minimum SINR required for reliable reception. In general, the value of \(\beta\) depends on the choice of coding and modulation. An outage occurs when the SINR falls below \(\beta\).

In \[14\], closed-form expressions were found for the outage probability conditioned on the particular network geometry and shadowing factors. Let \(\Omega = \{\Omega_0, \ldots, \Omega_M\}\) represent the set of normalized powers. Conditioning on \(\Omega\), the outage probability is

\[
\epsilon = P[\gamma \leq \beta | \Omega].
\]

(7)

Restricting the Nakagami parameter \(m_0\) of the channel between the reference transmitter and receiver to be integer-valued, the outage probability conditioned on \(\Omega\) is found in
is assumed that the summation in (10) is over all sets of indices that sum to $t$, and

$$G_{\ell} (\Psi_i) = \begin{cases} 1 - p_i (1 - \Psi_i^{m_i}) & \text{for } \ell = 0 \\ p_i \frac{\Gamma (\ell + m_i)}{\Gamma (m_i) \ell} \Psi_i^{m_i + \ell} & \text{for } \ell > 0. \end{cases}$$  (11)

The proof is found in Section III of [14].

V. EXAMPLES

In the examples, all distances are normalized to the network radius so that $r_{\text{net}} = 1$, and the reference receiver is located at the center of the network unless otherwise stated. As shown in Fig. 1, the reference transmitter is at coordinate $X_0 = 1/6$ and $M = 30$ potentially interfering mobiles are placed according to a uniform clustering model with a radius $r_{\text{ex}} = 1/12$ exclusion zone surrounding each mobile. This radius is small enough that the exclusion zones do not significantly impede the movements of mobiles. Once the mobile locations $\{X_i\}$ are realized, the $\{\Omega_i\}$ are determined by assuming a path-loss exponent $\alpha = 3.5$ and a common transmit power ($P_i / P_0 = 1$ for all $i$). Both unshadowed and shadowed environments are considered. Although the model permits nonidentical $G_i$, it is assumed that $G_i$ is a constant equal to $G_e$ (the effective processing gain) for all interference signals. Both spread and unspread systems are considered, with $G_e = 1$ for the unspread system and $G_e = 48$ for the spread system, corresponding to a typical direct-sequence waveform with $G = 32$ and $h(\tau_i) = 2/3$. Although the model permits nonidentical $p_i$ in the range $[0, 1]$, the value of $p_i = 0.5$ for all active $X_i$ is chosen as an example, corresponding to a half-duplex mobile terminal with a full input buffer and a symmetric data transmission rate to its peer terminal. The Nakagami parameters are $m_0 = 3$ for the desired signal and $m_i = 1$ for the interfering mobiles (i.e., Rayleigh fading). Using two different Nakagami factors (i.e., mixed fading) is justified by the fact that the reference transmitter is usually within the line-of-sight of the reference receiver while the interfering mobiles are not in the line-of-sight. Therefore, the interference signals are subject to more severe fading. The SINR threshold is $\beta = 0$ dB, which corresponds to the unconstrained AWGN capacity limit for a rate-1 channel code, assuming complex-valued inputs.

**Example #1.** Suppose that CSMA is not used, and therefore there is no guard zone beyond the fundamental exclusion zone; i.e., $r_g = r_{\text{ex}}$. All 30 potentially interfering mobiles in Fig. 1 remain active and contribute to the overall interference at the reference receiver. The outage probability for this network is shown in Fig. 2 by the two curves without markers, corresponding to the unspread ($G_e = 1$) and direct-sequence spread ($G_e = 48$) networks. Without direct-sequence spreading, the outage probability is quite high, for instance $\varepsilon = 0.4$ at $\Gamma = 20$ dB. Spreading reduces the outage probability by about three orders of magnitude at high SNR, although this comes at the cost of increased required bandwidth.

**Example #2.** The outage probability can be reduced by using CSMA to impose guard zones beyond the fundamental exclusion zone. Suppose that a guard zone of radius $r_g = 1/4$ is used, as shown in Fig. 1. Potentially interfering mobiles are deactivated according the following procedure, which is equivalent to Matern thinning. First, the reference transmitter $X_0$ is activated. Next, each potentially interfering mobile is...
considered in the order it was placed. For each mobile, a check is made to see if it is in the guard zone of a previously placed active mobile. Since mobiles are indexed according to the order of placement, $X_1$ is first considered for possible deactivation; if it lies in the guard zone of $X_0$, it is deactivated, and otherwise it is activated. The process repeats for each subsequent $X_i$, deactivating it if it falls within the guard zone of any active $X_j$, $j < i$, or otherwise activating it.

In Fig. 1 active mobiles are indicated by filled circles and deactivated mobiles are indicated by unfilled circles. The guard zone around each active mobile is indicated by a dashed circle of radius $r_g$. The reference receiver has not been assigned a CSMA guard zone, which reflects the fact that it has none while it is receiving the initial RTS. In the given example, 15 mobiles have been deactivated, and the remaining 15 mobiles remain active. The outage probability of the network with deactivated mobiles is shown in Fig. 2 for $G_a = \{1, 48\}$ by the two curves with markers. The performance of the unspread network improves dramatically with a guard zone, being reduced by two orders of magnitude at high SNR. The DS-CDMA network, which already had superior performance without CSMA, has improved performance, but the improvement only becomes significant at a high SNR.

VI. SPATIAL AVERAGING

Because it is conditioned on $\Omega$, the outage probability often varies significantly from one network realization to the next. The conditioning on $\Omega$ can be removed by averaging the conditional outage probability $\epsilon$ with respect to the network geometry. This averaging can be done analytically only under certain limitations [13]. For more general cases of interest, the outage probability can be estimated through Monte Carlo simulation by generating many different $\Omega$, computing the outage probability of each, and taking the numerical average.

Suppose that $N$ networks are generated, and let $\epsilon_i$ be the outage probability of the $i^{th}$ network, whose vector of normalized powers is expressed as $\Omega_i$. The average outage probability may be found by taking the average:

$$\bar{\epsilon} = \mathbb{E}(\epsilon) = \frac{1}{N} \sum_{i=1}^{N} \epsilon_i. \quad (12)$$

Generating each $\Omega_i$ involves not only placing the mobiles according to the uniform clustering model, but also realizing the shadowing and activating mobiles that lie within the CSMA guard zones of the active mobiles.

As an example, Fig. 3 shows the average outage probability computed over a set of $N = 10,000$ networks, each generated the same way as the Examples given in Section VII. For $r_g = r_{ex}$, each network was generated by placing $M = 30$ mobiles with exclusion zone $r_{ex} = 1/12$ according to a uniform clustering model. For $r_g = 1/4$, the same set of networks was used that was already generated with $r_{ex} = 1/12$, but in each network, potentially interfering mobiles were deactivated if they were in the guard zone of an active mobile. Log-normal shadowing with $\sigma_s = 8$ dB was applied, and all other system parameters have the same values as in Section V.

VII. PERFORMANCE ANALYSIS

This section uses the analytical tools developed in the previous sections to investigate the relationship among essential parameters, including the CSMA guard-zone radius, the exclusion-zone radius, the spreading factor, the number of potentially interfering mobiles, the reference-receiver location, and the transmission distance. Results are quantified by the spatially averaged outage probability, transmission capacity, and average latency.

Throughout this section, it is assumed that the SNR is $\Gamma = 10$ dB, all channels undergo mixed fading ($m_0 = 3$ and $m_i = 1, i \geq 1$) with log-normal shadowing ($\sigma_s = 8$ dB), the SINR threshold is $\beta = 0$ dB, and $p_i = 0.5$ for the active mobiles. The spatially averaged outage probability is computed by averaging over $N = 10,000$ network realizations. Unless otherwise specified, the path-loss exponent is set to $\alpha = 3.5$, and the reference receiver is at the center of the network.

A. Transmission Capacity

While the outage probability is improved with CSMA due to the deactivation of potentially interfering mobiles, the overall network becomes less efficient due to the suppression of transmissions. The network efficiency can be quantified by the transmission capacity (TC) [5]:

$$\tau = (1 - \bar{\epsilon}) \lambda b \quad (13)$$

where $\lambda$ is the network density, which is the number of active mobiles per unit area, and $b$ is the link throughput in the absence of an outage, in units of bps. The TC represents the network throughput per unit area. Increasing the size of the guard zone generally reduces TC due to fewer simultaneous transmissions.

For a given value of $M$, the density without a CSMA guard zone remains fixed since all mobiles remain active. However, with a CSMA guard zone, the number of potentially interfering mobiles is random with a value that depends on the value of $r_g$, the locations of the mobiles, and their order of
placement, which affects how they are deactivated. As with outage probability, Monte Carlo simulation can be used to estimate TC.

B. Latency

Assume that a selective-repeat hybrid ARQ protocol is used. For such a protocol, the receiver attempts to decode the received packet, and signals back a positive acknowledgement (ACK) if it is correctly decoded or a negative acknowledgement (NACK) if it is not. By setting \( \beta \) to the minimum SINR required to correctly decode the packet, a retransmission will be required with probability \( \epsilon \). Assuming that type-I hybrid ARQ is used, the receiver flushes its memory and will attempt to independently decode each retransmitted packet.

Let \( T_s \) be the ARQ time-slot duration, and define latency to be the average time from when a packet is sent to when it is first successfully decoded. Let \( N_{ARQ} \) be the number of slots between the start of the initial transmission and the start of the retransmission, so that if a packet sent in time slot \( n \) fails, then it is retransmitted in slot \( n + N_{ARQ} \). For such a protocol, the average probability of success in the first slot is \((1 - \bar{\epsilon})\), and the average probability of success in slot \( 1 + kN_{ARQ} \) is \( \bar{\epsilon}^k(1 - \bar{\epsilon}) \). Taking the expected value of the number of slots before success and multiplying by the slot duration, the average latency \( D \) is found to be

\[
D = T_s \left( \frac{N_{ARQ}}{1 - \bar{\epsilon}} - (1 - \bar{\epsilon})(N_{ARQ} - 1) \right). \tag{14}
\]

C. Effect of \( r_{ex} \) and \( r_g \)

To investigate the influence of the exclusion-zone radius \( r_{ex} \) and guard-zone radius \( r_g \) on the network performance, the outage probability, transmission capacity, and latency were determined over a range of \( r_{ex} \) and \( r_g \). To remove the dependence on the transmitter-receiver separation, the guard and exclusion zones are normalized with respect to \( ||X_0|| \). The normalized guard-zone radius was varied over \( 0.5 \leq r_g/||X_0|| \leq 3 \), and three representative values of the normalized exclusion-zone radius were selected: \( r_{ex}/||X_0|| = \{1/4, 1/2, 3/4\} \). Both unspread and spread \((G_e = 48)\) networks were considered. The number of potentially interfering mobiles was set to \( M = 30 \) and \( r_{net}/||X_0|| = 6 \).

Fig. 4 shows the spatially averaged outage probability. For the unspread system, the outage probability is insensitive to the exclusion-zone radius but very sensitive to the guard-zone radius. This underscores the importance of a guard zone for an unspread network. Except when the guard-zone radius is relatively small, the outage probability of the spread system is insensitive to both the exclusion-zone and guard-zone radii.

Fig. 5 shows the transmission capacity as a function of the guard-zone radius. Although the outage probability at the reference receiver in an unspread network is insensitive to \( r_{ex} \), Fig. 5 shows that the TC is sensitive to \( r_{ex} \), especially at low \( r_g \). As \( r_{ex} \) increases, there are fewer nearby interfering mobiles that get deactivated by the guard zone. Thus, more mobiles remain active, and the transmission capacity increases even though the outage probability remains fixed. A similar behavior is seen for the spread network, which for small \( r_g \) has a significantly higher TC than the unspread network due to the lower outage probability. The main limitation to increasing \( r_{ex} \) is that it must be small enough that there is no significant impediment to the movements of mobiles. For both the spread and unspread network, the TC diminishes quickly with increasing \( r_g \), and at high \( r_g \), the TC is insensitive to the spreading factor and exclusion-zone radius.

The average latency is shown in Fig. 6 for \( N_{ARQ} = 6 \). As the latency is a function of the outage probability, the general trends are similar to those already observed in Fig. 4. In particular, the unspread network is more sensitive than the spread network to the choice of exclusion-zone radius, and the latency performance of the spread network is better than that of the unspread network. Except for small values of guard-zone radius, the latency is relatively insensitive to the radius of the exclusion zone. As the exclusion zone radius is limited in practice by the mobility constraints of the mobiles,
it follows that mobility does not have a significant effect on latency.

D. Effect of the Reference-Receiver Location

In a finite network, the average outage probability depends on the location of the reference receiver. Rather than leaving the reference receiver at the origin, Table II explores the change in outage probability when the reference receiver moves from the center of the network to its perimeter. The guard and exclusion zones are again normalized with respect to $||X_0||$. Two representative values for each of the normalized exclusion-zone and guard-zone radii are considered: $r_{ex}/||X_0|| = \{0, 1/2\}$ and $r_{g}/||X_0|| = \{1/2, 3/2\}$. Two different values of path-loss exponent are considered: $\alpha = 3$ and $\alpha = 4$. Both unspread and spread $(G_c = 48)$ networks are considered. The number of potentially interfering mobiles are set to $M = 30$, and $r_{net}/||X_0|| = 6$. When the reference receiver is located on the network perimeter, the reference transmitter is placed in the direction toward the network center.

Table II indicates that the average outage probability on the finite network’s perimeter $\bar{\epsilon}_p$ is considerably less than the average outage probability at the network’s center $\bar{\epsilon}_c$. This result cannot be predicted by the traditional infinite-network model, which cannot differentiate between $\bar{\epsilon}_c$ and $\bar{\epsilon}_p$. The reduction in outage probability is more significant for the unspread network and is less pronounced with increasing $G_c$. Both direct-sequence spreading and CSMA make the average performance less dependent on the location of the reference receiver inside the network.

E. Effect of the Number of Potentially Interfering Mobiles

In the previous figures, a fixed number of potentially interfering mobiles ($M = 30$) was placed prior to CSMA deactivation. In Fig. 7 $M$ is varied from 2 to 60, and the transmission capacity is shown for each value of $M$ for a normalized exclusion-zone radius $r_{ex}/||X_0|| = 1/2$, $r_{net}/||X_0|| = 6$, and both unspread and spread ($G_c = 48$) networks. Results are shown for a CSMA guard zone with normalized radius $r_g/||X_0|| = 3/2$ and for a system that uses no additional CSMA guard zone; i.e., $r_g/||X_0|| = r_{ex}/||X_0|| = 1/2$.

As observed previously for just $M = 30$, the transmission capacity of the DS-CDMA network is higher than that of the unspread network, and using a CSMA guard zone reduces transmission capacity. The transmission capacity of the DS-CDMA network increases roughly linearly with $M$ in the absence of a guard zone, but the increase is sublinear for the unspread network or when a guard zone is used. At $M = 60$, the transmission capacity of the unspread network is approximately the same with and without a guard zone.

F. Effect of Transmitter Distance

The previous examples assume that the distance between the reference transmitter and receiver is fixed with respect to the network radius. However, performance will depend on this distance. Fig. 8 shows $\tau/b$, the normalized TC as a function
of the distance $||X_0||$ between the reference transmitter and receiver. All distances are normalized to the network radius so that $r_{\text{net}} = 1$. The exclusion-zone radius is set to $r_{\text{ex}} = 1/12$, and both unspread and spread ($G_e = 48$) networks are considered. Results are shown for a CSMA guard zone with radius $r_g = 1/4$ and for a system that uses no additional CSMA guard zone.

It is observed that increasing the transmission distance reduces the TC due to the increase in the number of interfering mobiles that are closer to the receiver than the reference transmitter, but this reduction is more gradual with the spread system than the unspread one. As $||X_0||$ increases, an increased guard zone alleviates the potential near-far problems. Consequently, the rate of TC degradation is made more gradual by the use of CSMA, and at sufficiently large transmitter distances, a system with CSMA outperforms a system without it.

The CSMA guard zone decreases transmission capacity at short transmission distances, but increases it at long distance. This is because at a long transmission distance, the received signal power from the far reference transmitter will be weak, while the received powers from the nearby interferers will be relatively high. To overcome this near-far problem, nearby interferers need to be deactivated in order for the SINR threshold to be met. However, at short distances, the signal power from the nearby reference transmitter will already be strong enough, and deactivating the interferers is unnecessary and harmful to the transmission capacity due to a reduction of simultaneous transmissions.

$G$. Tradeoff Between Spreading Factor and $r_g$

The average outage probability depends on both the effective spreading factor $G_e$ and guard-zone radius $r_g$. Fig. 9 explores the tradeoff between these two parameters by plotting the minimum $r_g$ required to achieve transmission capacity $\tau \geq 15b$ for a given value of $G_e$. As with Fig. 8 all distances are normalized to the network radius, $r_{\text{ex}} = 1/12$, the reference transmitter is placed at one of four distances, and two values of $M$ are considered. The curves are generated using a Monte Carlo approach as follows. For each value of $M$, $N = 10,000$ networks are realized by drawing from a uniform clustering process. For each $r_g$, nodes are deactivated using the Matern thinning policy described in Example #2, and new values of $G_e$ are repeatedly picked until one is found that has a corresponding $\bar{\epsilon}$ within the range $0.0999 \leq \bar{\epsilon} \leq 0.1010$. It is observed that the required spreading factor increases as the guard-zone radius decreases until the guard zone coincides with the exclusion zone.

As with the outage probability, the transmission capacity depends on both the effective spreading factor $G_e$ and guard-zone radius $r_g$. Fig. 10 explores the tradeoff between these two parameters by plotting the minimum $G_e$ required to achieve transmission capacity $\tau \geq 15b$ for a given value of $r_g$. As with Fig. 9 all distances are normalized to the network radius, $r_{\text{ex}} = 1/12$, the reference transmitter is placed at one of four distances, and two values of $M$ are considered. The curves are generated using a Monte Carlo approach similar to the one used to generate Fig. 9. The need for a larger $G_e$ as $r_g$ grows is consistent with the observation that the transmission capacity generally decreases as $r_g$ (see Fig. 5). Although the average outage probability at the reference receiver decreases with increasing $r_g$, the number of interfering mobiles is reduced due to the thinning, and hence the transmission capacity is reduced.

VIII. CONCLUSIONS

The analysis developed in this paper allows the tradeoffs between exclusion zones and CSMA guard zones to be explored for DS-CDMA and unspread networks. The spreading factor and the guard-zone size provide design flexibility in achieving specified levels of average outage probability and transmission capacity. The advantage of an exclusion zone over a CSMA guard zone is that since the network is not thinned, the number
of active mobiles remains constant, and higher transmission capacities can be achieved. However, practical constraints of mobility and positioning of mobile terminals can sometimes prevent the use of a significant exclusion zone. In these cases, a CSMA guard zone might be the only option. While exclusion zones are advantageous when they can be imposed, when they cannot, a CSMA guard zone still provides performance benefits.

The computational results indicate a number of general tendencies for networks with exclusion and guard zones. Increases in the size of the exclusion zone are beneficial, but there are practical limits to this size, and the transmission capacity associated with larger exclusion zones diminishes quickly with increasing guard zones. The CSMA guard zone is much more useful in networks with unspread systems than those with direct-sequence spreading. Both spreading and a CSMA guard zone make the outage probability less dependent on the location of the receiver within the network. The spreading factor required for a specified outage probability less increases with decreases in the guard zone until the guard zone coincides with the exclusion zone. The rate of transmission-capacity degradation as the transmitter distance increases is made more gradual by the use of CSMA, and at sufficiently large transmitter distances, a system with CSMA outperforms a system without it. Although both the average outage probability and the latency of a receiver decrease with increasing guard zones, the number of interfering mobiles is reduced due to the thinning, and hence the transmission capacity is reduced.

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