Multipath TCP: Analysis, Design and Implementation

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Abstract—Multi-path TCP (MP-TCP) has the potential to greatly improve application performance by using multiple paths transparently. We propose a fluid model for a large class of MP-TCP algorithms and identify design criteria that guarantee the existence, uniqueness, and stability of system equilibrium. We clarify how algorithm parameters impact TCP-friendliness, responsiveness, and window oscillation and demonstrate an inevitable tradeoff among these properties. We discuss the implications of these properties on the behavior of existing algorithms and motivate a new design that generalizes existing algorithms and strikes a good balance among TCP-friendliness, responsiveness, and window oscillation. We have implemented our algorithm in the Linux kernel. We use our prototype to compare the new algorithm with existing MP-TCP algorithms.

I. INTRODUCTION

Traditional TCP uses a single path through the network even though multiple paths are usually available in today’s communication infrastructure; e.g., most smartphones are enabled with both cellular and WiFi access, and servers in data centers are connected to multiple routers. Multi-path TCP (MP-TCP) has the potential to greatly improve application performance by using multiple paths transparently. It is being standardized by the MP-TCP Working Group of the Internet Engineering Task Force (IETF) [2]. In this paper we present a fluid model of MP-TCP and study how protocol parameters affect structural properties such as the existence, uniqueness and stability of equilibrium, the tradeoffs among TCP friendliness, responsiveness and window oscillation. These properties motivate a new algorithm that generalizes existing MP-TCP algorithms.

Various congestion control algorithms have been proposed as an extension of TCP NewReno for MP-TCP. A straightforward extension is to run TCP NewReno on each subpath, e.g. [3], [4]. This algorithm however can be highly unfriendly when it shares a path with a single-path TCP user. This motivates the Coupled algorithm which is fair because it has the same underlying utility function as TCP NewReno, e.g. [5], [6]. It is found in [7], however that the Coupled algorithm respond slowly in a dynamic network environment. A different algorithm is proposed in [7] (which we refer to as the Max algorithm) which is more responsive than the Coupled algorithm and still reasonably friendly to single-path TCP users. Recently, opportunistic linked increase algorithm (OLIA) is proposed as a variant of Coupled algorithm that is as friendly as the Coupled algorithm but more responsive [8]. See [9] for more references to early work on multi-path congestion control.

Our goal is to develop structural understanding of MP-TCP algorithms so that we can systematically tradeoff different properties such as TCP friendliness, responsiveness, and window oscillation that can be detrimental to applications that require a steady throughput. For single-path TCP, one can associate a strictly concave utility function with each source so that the congestion control algorithm implicitly solves a network utility maximization problem [9]–[11]. The convexity of this underlying utility maximization guarantees the existence, uniqueness, and stability of most single-path TCP algorithms.

For many MP-TCP proposals considered by IETF, it will be shown that the utility maximization interpretation fails to hold in general, necessitating the need for a different approach to understanding the equilibrium properties of these algorithms. Moreover the relations among different performance metrics, such as fairness, responsiveness and window oscillation, need to be clarified.

The main contributions of this paper are three-fold. First we present a fluid model that covers a broad class of MP-TCP algorithms and identify the exact property that allows an algorithm to have an underlying utility function. This implies that some MP-TCP algorithms, e.g., the Max algorithm [7], has no associated utility function. We prove conditions on protocol parameters that guarantee the existence and uniqueness of the equilibrium, and its asymptotical stability. Indeed algorithms that fail to satisfy these conditions, e.g. the Coupled algorithm, can be unstable and can have multiple equilibria as shown in [7]. Second we clarify how protocol parameters impact TCP friendliness, responsiveness, and window oscillation and demonstrate the inevitable tradeoff among these properties. Finally, based on our understanding of the design space, we propose a new MP-TCP algorithm that generalizes existing algorithms and strikes a good balance among these properties. This algorithm has been implemented in the Linux kernel and we evaluate its performance using our Linux prototype.

We now summarize our proposed MP-TCP algorithm. Each source $s$ has a set of routes $r$. Each route $r$ maintains a congestion window $w_r$ and measures its round-trip time $\tau_r$. The window adaptation is as follows:

- For each ACK on route $r \in s$,

$$w_r \leftarrow w_r + \frac{x_r}{\tau_r (\sum x_k)} \left( \frac{1 + \alpha_r}{2} \right) \left( \frac{4 + \alpha_r}{5} \right)$$

(1)
II. MULTIPATH TCP MODEL

In this section we first propose a fluid model of MP-TCP and then use it to model MP-TCP algorithms in the literature. Unless otherwise specified, a boldface letter $x \in \mathbb{R}^n$ denotes a vector with components $x_i$. We use $x_{-i} := (x_1, \ldots, x_{i-1}, x_{i+1}, \ldots, x_n)$ to denote the $n-1$ dimensional vector without $x_i$, and $\|x\|_k := (\sum x_i^k)^{1/k}$ to denote the $L_k$-norm of $x$. Given two vectors $x, y \in \mathbb{R}^n$, $x \geq y$ means $x_i \geq y_i$ for all components $i$. A capital letter denotes a matrix or a set, depending on the context. A symmetric matrix $P$ is said to be positive (negative) semidefinite if $x^T P x \geq 0$ (respectively, $x^T P x \leq 0$) for any $x$, and positive (negative) definite if $x^T P x > 0$ (respectively, $x^T P x < 0$) for any $x \neq 0$. For any matrix $P$, define $[P]^+ := (P + P^T)/2$ to be its symmetric part. Given two arbitrary matrices $A$ and $B$ (not necessarily symmetric), $A \geq B$ means $[A - B]^+$ is positive semidefinite. For a vector $x$, $\text{diag}(x)$ is a diagonal matrix with entries given by $x$.

A. Fluid model

Consider a network that consists of a set $L = \{1, \ldots, |L|\}$ of links with finite capacities $c_l$. The network is shared by a set $S = \{1, \ldots, |S|\}$ of sources. Available to source $s \in S$ is a fixed collection of routes (or paths) $r$. A route $r$ consists of a set of links $l$. We abuse notation and use $s$ both to denote a source and the set of routes $r$ available to it, depending on the context. Likewise, $r$ is used both to denote a route and the set of links $l$ in the route. Let $R := \{r \mid r \in S, s \in S\}$ be the collection of all routes. Let $H \in \{0,1\}^{|L| \times |R|}$ be the routing matrix: $H_{lr} = 1$ if link $l$ is in route $r$ (denoted by ‘$l \in r$’), and 0 otherwise.

For each route $r \in R$, $\tau_r$ denotes its round trip time (RTT). For simplicity we assume $\tau_r$ are constants. Each source $s$ maintains a congestion window $w_r(t)$ at time $t$ for every route $r \in s$. Let $x_r(t) := w_r(t)/\tau_r$ represent the sending rate on route $r$. Each link $l$ maintains a congestion price $p_l(t)$ at time $t$. Let $q_r(t) := \sum_{l \in L} H_{lr} p_l(t)$ be the aggregate price on route $r$. In this paper $p_l(t)$ represents the packet loss probability at link $l$ and $q_r(t)$ represents the approximate packet loss probability on route $r$.

We associate three state variables $(x_r(t), w_r(t), q_r(t))$ for each route $r \in s$. Let $x_s(t) := (x_r(t), r \in s)$, $w_s(t) := (w_r(t), r \in s)$, $q_s(t) := (q_r(t), r \in s)$. Then $(x_s(t), w_s(t), q_s(t))$ represents the corresponding state variables for each source $s \in S$. For each link $l$, let $y_l(t) := \sum_{r \in R} H_{lr} x_r(t)$ be its aggregate traffic rate.

Congestion control is a distributed algorithm that adapts $x(t)$ and $p(t)$ in a closed loop. Motivated by the AIMD algorithm of TCP Newreno, we model MP-TCP by

$$\dot{x}_r = k_r(x_r) \cdot (\hat{\phi}(x_s) - q_r)^+ \quad r \in s \in S$$

$$\dot{y}_l = \gamma_l (y_l - c_l)^+ \quad l \in L,$$

where $(a)^+ = a$ for $x > 0$ and $\max\{0, a\}$ for $x \leq 0$. We omit the time $t$ in the expression for simplicity. We model how sending rates are adapted in the congestion avoidance phase of TCP at each end system and models how the congestion price is (often implicitly) updated at each link. The MP-TCP algorithm installed at source $s$ is specified by $(K_s, \Phi_s)$, where $K_s(x_s) := (k_r(x_s), r \in s)$ and $\Phi_s(x_s) := (\hat{\phi}(x_s), r \in s)$. Here $K_s(x_s)$ is a vector of positive gains that determines the dynamic property of the algorithm. $\Phi_s(x_s)$ determines the equilibrium properties of the algorithm. The link algorithm is specified by $\gamma_l$, where $\gamma_l > 0$ is a positive gain that determines the dynamic property. This is a simplified model for the RED algorithm that assumes the loss probability is proportional to the backlog, and is used in, e.g., [10], [11].

B. Existing MP-TCP algorithms

We first show how to relate the fluid model to the window-based MP-TCP algorithms proposed in the literature. On each route $r$ the source increases its window at the return of each ACK. Let this increment be denoted by $I_r(w_s)$ where $w_s$ is the vector of window sizes on different routes of source $s$. The source decreases the window on route $r$ when it sees a packet loss on route $r$. Let this decrement be denoted by $D_r(w_s)$. Then most loss based MP-TCP algorithms take the form of the following pseudo code:

- For each packet loss on route $r$, $w_r \leftarrow w_r + I_r(w_s)$.
- For each ACK on route $r$, $w_r \leftarrow w_r - D_r(w_s)$.

We now model the above pseudo codes by the fluid model [5]. Let $\delta w_r$ be the net change to window on route $r$ in each round trip time. Then $\delta w_r$ is roughly

$$\delta w_r = (I_r(w_s)(1 - q_r) - D_r(w_s) q_r) w_r$$

$$\approx (I_r(w_s) - D_r(w_s) q_r) w_r$$

since the loss probability $q_r$ is small. On the other hand

$$\delta w_r \approx \dot{x}_r \tau_r$$

Hence

$$\dot{x}_r = \frac{\tau_r}{\tau_r} (I_r(w_s) - D_r(w_s) q_r)$$

From (3) we have

$$\begin{align*}
k_r(x_s) &= \frac{x_s}{\tau_r} D_r(w_s) \\
\hat{\phi}(x_s) &= \frac{1}{\tau_r} D_r(w_s)
\end{align*}$$

(5)
We now apply this to the algorithms in the literature. We first summarize these algorithms in the form of a pseudo-code and then use\(^5\) to derive parameters \(k_r(x_s)\) and \(\phi_r(x_s)\) of the fluid model\(^5\).

Single-path TCP (TCP-NewReno): Single-path TCP is a special case of MP-TCP algorithm with \(|s| = 1\). Hence \(x_s\) is a scalar and we identify each source with its route \(r = s\). TCP-NewReno adjusts the window as follows:

- For each ACK on route \(r\), \(w_r \leftarrow w_r + 1/w_r\).
- For each loss on route \(r\), \(w_r \leftarrow w_r/2\).

From\(^5\), this can be modeled by the fluid model\(^5\) with

\[
k_r(x_s) = \frac{1}{2}x_s^2, \quad \phi_r(x_s) = \frac{2}{r^2x_s^2}
\]

We now summarize some existing MP-TCP algorithms, all of which degenerate to TCP NewReno if there is only one route per source.

EWTCP\(^7\): EWTCP algorithm applies TCP-NewReno like algorithm on each route independently of other routes. It adjusts the window on multiple routes as follows:

- For each ACK on route \(r\), \(w_r \leftarrow w_r + a/w_r\).
- For each loss on route \(r\), \(w_r \leftarrow w_r/2\).

From\(^5\), this can be modeled by the fluid model\(^5\) with

\[
k_r(x_s) = \frac{1}{2}x_s^2, \quad \phi_r(x_s) = \frac{2a}{r^2x_s^2}
\]

where \(a > 0\) is a constant.

Coupled MPTCP\(^3, 6\): The Coupled MPTCP algorithm adjusts the window on multiple routes in a coordinated fashion as follows:

- For each ACK on route \(r\), \(w_r \leftarrow w_r + \frac{w_r/\tau_r^2}{(\sum_{k \in s} w_k/\tau_k)^2}\).
- For each loss on route \(r\), \(w_r \leftarrow w_r/2\).

From\(^5\), this can be modeled by the fluid model\(^5\) with

\[
k_r(x_s) = \frac{1}{2}x_s^2, \quad \phi_r(x_s) = \frac{2}{r^2(x_s/x_r)^2}
\]

Semicoupled MPTCP\(^7\): The Semicoupled MPTCP algorithm adjusts the window on multiple routes as follows:

- For each ACK on route \(r\), \(w_r \leftarrow w_r + \frac{1}{r^2(\sum_{k \in s} w_k/\tau_k)^2}\).
- For each loss on route \(r\), \(w_r \leftarrow w_r/2\).

From\(^5\), this can be modeled by the fluid model\(^5\) with

\[
k_r(x_s) = \frac{1}{2}x_s^2, \quad \phi_r(x_s) = \frac{2}{x_r^2(r^2(\sum_{k \in s} x_k))^2}
\]

Max MPTCP\(^7\): The Max MPTCP algorithm adjusts the window on multiple routes as follows:

- For each ACK on route \(r\), \(w_r \leftarrow \min \left\{ \frac{\max(w_k/\tau_k)}{(\sum_{k \in s} w_k/\tau_k)^2}, \frac{1}{w_r} \right\}\).
- For each loss on route \(r\), \(w_r \leftarrow w_r/2\).

From\(^5\), this can be modeled by the fluid model\(^5\) with

\[
k_r(x_s) = \frac{1}{2}x_s^2, \quad \phi_r(x_s) = \frac{2\max(x_k/\tau_k)}{x_r^2(\sum_{k \in s} x_k)^2}
\]

where we have ignored taking the minimum with the \(1/w_r\) term since the performance is mainly captured by \(\sum_{k \in s} x_k\).

Recently, OLIA MP-TCP algorithm\(^8\) is shown to achieve good performance in many scenarios. OLIA uses complicated feedback congestion control signals and cannot be modeled by\(^5\). We do, however, include OLIA in our Linux-based performance evaluation in section\(^5\).

III. STRUCTURAL PROPERTIES

Throughout this paper we assume, for all \(x_s, r \in s, s \in S\), \(k_r(x_s) > 0\) and \(\phi_r(x_s) = 0\) only if \(x_k = \infty\) for some \(k \in s\). A point \((x,p)\) is called an equilibrium of\(^3\) if it satisfies, for all \(r \in s, s \in S\) and \(l \in L\),

\[
k_r(x_s) \phi_r(x_s) = 0 = \gamma_l (y_l - c_l/p_l)
\]

or equivalently,

\[
x_r \geq 0, \quad \phi_r(x_s) \leq q_r \quad \text{and} \quad \phi_r(x_s) = q_r \quad \text{if} \quad x_r > 0
\]

We make two remarks. First an equilibrium \((x,p)\) does not depend on \(K_s\), but only on \(\Phi_s\). The design \((K_s,s \in S)\) however affects dynamic properties such as stability and responsiveness as we show below. Second, since \(k_r(x_s) > 0\) and \(\phi_r(x_s) = 0\) only if \(x_k = \infty\) for some \(k \in s\) by assumption, any finite equilibrium \((x,p)\) must have \(q_r > 0\) for all \(r\). In the following we always restrict ourselves to finite equilibria.

In this section we denote an MP-TCP algorithm by \((K,\Phi) := (K_s,\Phi_s, s \in S)\). We characterize MP-TCP designs \((K,\Phi)\) that guarantee the existence, uniqueness, and stability of system equilibrium. We identify design criteria that determine TCP-friendliness, responsiveness and window oscillation and prove an inevitable tradeoff among these properties. We discuss in the next section the implications of these structural properties on existing algorithms. All proofs are relegated to the Appendices.

A. Summary

We first present some properties of an MP-TCP algorithm \((K,\Phi)\) that we have identified. We then interpret them and summarize their implications.

C0: For each \(s \in S\) and each \(x_s\), the Jacobians of \(\Phi_s(x_s)\) is continuous and symmetric, i.e.,

\[
\frac{\partial \Phi_s}{\partial x_s}(x_s) = \left[ \frac{\partial \Phi_s}{\partial x_s}(x_s) \right]^T
\]

C1: For each \(s \in S\) there exists a nonnegative solution \(x_s := x_s(p)\) to\(^6\) for any finite \(p \geq 0\) such that \(q_r > 0\) for all

| Algorithm | C0 | C1 | C2, C3 | C4 | C5 |
|-----------|----|----|--------|----|----|
| EWTCP     | Yes | Yes | Yes    | Yes | Yes |
| Coupled   | Yes | Yes | No     | Yes | Yes |
| Semicoupled | No  | Yes | Yes    | Yes | Yes |
| Max       | No  | Yes | Yes    | Yes | Yes |
| Generalized | No  | Yes | Yes    | Yes | Yes |

Theorem 3.1 [3.2, 3.3, 3.5, 3.4, 3.6]
C4: For each r, \( y^*_r(p) := \sum_{s \in S} H_{1r} x_r(p) \) is the aggregate traffic at link l from source s.

C2: For each s \( \in S \) and each \( x_s \), \( \Phi_s(x_s) \) is differentiable with respect to \( x_s \); moreover, the symmetric part \( \frac{\partial^2 \Phi_s(x_s)}{\partial x_s \partial x_s} \) is negative definite.

C3: For each \( r \in R \), \( \phi_r(x_s) \) is the route's queueing delay.

C4: For each \( r \in R \), s \( \in S \), \( D_s(x_r) := \left[ \frac{\partial x_r}{\partial y^*_r} \right]^{-1} \).

C5: For each \( r \in R \) and each \( x_{-r} \), \( \lim_{x_r \to \infty} \phi_r(x_r) = 0 \).

These design criteria are intuitive and usually (but not always) satisfied; see Table I.

Conditions C1–C3 guarantee the existence, uniqueness, and global asymptotic stability of the equilibrium (x, p) (Theorems 3.2 and 3.3). C1 says that the aggregate traffic rate through a link l from source s decreases when the congestion price \( p_l \) increases. C2 implies that at steady state \( x_s \), \( q_s \) is perturbed by \( \delta x_s \), \( \delta q_s \) respectively, then \( (\delta x_s)^T \delta q_s < 0 \).

C3 says that in the case of single-path TCP (SP-TCP), one can associate a utility function \( U_s(x_s) \) with each flow \( x_s \). The existence of \( D_s(x_r) \) is optimal for the associated dual problem. Here \( \phi_r(x_s) \) is a function of \( (x, p) \) and \( \phi_r(x_s) \) is only a function of \( (x, p) \) but not \( (x, p) \).

B. Utility maximization

For single-path TCP (SP-TCP), one can associate a utility function \( U_s(x_s) \in \mathbb{R}_+ \rightarrow \mathbb{R} \) with each flow \( x_s \) (\( x_s \) is a scalar and \(|s| = 1 \)) and interpret (3)–(4) as a distributed algorithm to maximize aggregate users' utility, e.g. [9]–[12]. Indeed, for SP-TCP, an (x, p) is an equilibrium if and only if x is optimal for

\[
\text{maximize } \sum_{s \in S} U_s(x_s) \text{ s.t. } y_l \leq c_l, \forall l \in L
\]

where \( x_s := (x_r, r \in s) \) is the rate vector of flow s and \( U_s: \mathbb{R}^{|s|} \rightarrow \mathbb{R} \) is a concave function.

Theorem 3.1 (utility maximization): There exists a twice continuously differentiable and concave utility function \( U_s(x_s) \) such that an equilibrium (x, p) of (3)–(4) solves (9) and its dual problem if and only if condition C0 holds.

C. Existence, uniqueness and stability of equilibrium

Even though a multipath TCP algorithm \((K, \Phi)\) may not have a utility maximization interpretation, a unique equilibrium exists if conditions C1–C3 are satisfied.

Theorem 3.2 (existence and uniqueness):

1) Suppose C1 holds. Then (3)–(4) has at least one equilibrium.

2) Suppose C2 and C3 hold. Then (3)–(4) has at most one equilibrium.
Theorem 3.3 (stability): Suppose C1–C3 hold and \( k_r(x_s) \equiv k_r(x_r) \) for all \( r \in R \). Then the unique equilibrium \((x^*, p^*)\) is globally asymptotically stable. In particular, starting from any initial point \( x(0) \in \mathbb{R}^{|R|} \) and \( p(0) \in \mathbb{R}^{|L|} \), the trajectory \((x(t), p(t))\) generated by the MP-TCP algorithm \((3)\)–\((4)\) converges to the equilibrium \((x^*, p^*)\) as \( t \to \infty \).

Our proposed algorithm does not satisfy \( k_r(x_s) \equiv k_r(x_r) \) even though it seems to be stable in our experiments. This condition is only sufficient and needed in our Lyapunov stability proof; see Appendix C. When \( k_r(x_s) \) depends on \( x_s \), one can replace \( k_r(x_r) \) in the definition of the Lyapunov function \( V \) in \((21)\) with \( k_r(x_s) \) evaluated at the equilibrium and the same argument there proves that \((x^*, p^*)\) is (locally) asymptotically stable. Also see Theorem 3.5 below for an alternative proof of local stability.

D. TCP friendliness

Informally, an MP-TCP flow is said to be ‘TCP friendly’ if it does not dominate the available bandwidth when it shares the same network with a SP-TCP flow \([2]\). To define this precisely we use the test network shared by a SP-TCP flow and a MP-TCP flow under test as shown in Fig. 1.

All paths traverse a single bottleneck link with capacity \( c \), with all other links with capacities strictly higher than \( c \). The links have fixed but possibly different delays. To compare the friendliness of two MP-TCP algorithms \( \hat{M} \) \(=\) \( (K, \hat{\Phi}) \) and \( M \) \(=\) \( (\hat{K}, \hat{\Phi}) \), suppose that \( \hat{M} \) shares the test network with a SP-TCP it achieves a throughput of \( \|x\|_1 \) in equilibrium aggregated over the available paths (the SP-TCP therefore attains a throughput of \( c - \|x\|_1 \)). Suppose \( \hat{M} \) achieves a throughput of \( \|x\|_1 \) in equilibrium when it shares the test network with the same SP-TCP. Then we say that \( \hat{M} \) is friendlier (or more TCP-friendly) than \( M \) if \( \|x\|_1 \leq \|x\|_1 \), i.e., if \( \hat{M} \) receives no more bandwidth than \( M \) does when they separately share the test network in Fig. 1 with the same SP-TCP flow.

From the theory for single-path TCP \((|s| = 1 \text{ for all } s \in S)\), it is known that a design is more TCP-friendly if it has a smaller marginal utility \( U'_s(x_s) = \Phi_s(x_s) \). The same intuition holds for MP-TCP algorithms even though the utility functions may not exist for MP-TCP algorithm.

Theorem 3.4 (friendliness): Consider two MP-TCP algorithms \( \hat{M} \) \(=\) \( (K, \hat{\Phi}) \) and \( M \) \(=\) \( (\hat{K}, \hat{\Phi}) \). Suppose both satisfy C1–C4. Then \( \hat{M} \) is friendlier than \( M \) if \( \hat{\Phi}_s(x_s) \leq \hat{\Phi}_s(x_s) \) for all \( s \in S \).

E. Responsiveness around equilibrium

Suppose conditions C1–C3 hold and there is a unique equilibrium \( z^* := (x^*, p^*) \). Assume all links in \( L \) are active with \( p^*_j > 0 \); otherwise remove from \( L \) all links with prices \( p^*_j = 0 \). Let \( \delta z(t) := z(t) - z^* \). The behavior of \((3)\)–\((4)\) around the equilibrium is defined by the linearized system:

\[
\delta \dot{z} = J^* \delta z(t)
\]

Here \( J^* \) is the Jacobian of \((3)\)–\((4)\) at the equilibrium \( z^* \):

\[
J^* := J(x^*) := \begin{bmatrix} \Lambda_k \frac{\partial \Phi}{\partial x} - \Lambda_k^T \Lambda_y H \\
\Lambda_y H & 0 \end{bmatrix}
\]

where \( \Lambda_k = \text{diag} \{ k_r(x^*_r), r \in R \}, \Lambda_y = \text{diag} \{ \gamma_{il}, l \in L \} \), and \( \frac{\partial \Phi}{\partial x} \) is evaluated at \( x^* \).

The stability and responsiveness of the linearized system \((10)\) (how fast does the system converges to the equilibrium locally) is determined by the real parts of the eigenvalues of \( J^* \). Specifically the linearized system is stable if the real parts of all eigenvalues of \( J^* \) are negative; moreover the more negative the real parts are the faster the linearized system converges to the equilibrium. We now show that the linearized system \((10)\) is stable (i.e., converges exponentially fast to \( z^* \) locally) and characterize its responsiveness in terms of the design choices \((K, \Phi)\).

Let \( Z = \{ z := (x, p) \in \mathbb{C}^{(|R|+|L|)} \mid \|z\|_2 = 1 \} \).

Theorem 3.5 (responsiveness): Suppose C1–C3 hold. Then

1) The linearized system \((10)\) is stable, i.e., \( \text{Re}(\lambda) < 0 \) for any eigenvalue \( \lambda \) of \( J^* \). Moreover \( \text{Re}(\lambda) \leq \bar{\lambda}(J^*) \) where

\[
\bar{\lambda}(J^*) := \max_{z \in Z} \left\{ \frac{x^H \left[ \frac{\partial \Phi}{\partial x} \right]^+ x}{x^H \Lambda_k^{-1} x + p^H \Lambda_y^{-1} p} \right\} \leq 0
\]

where \( \Lambda_k \) and \( \frac{\partial \Phi}{\partial x} \) are evaluated at the equilibrium point \( z^* \).

2) For two MP-TCP algorithms \((K, \hat{\Phi})\) and \((\hat{K}, \hat{\Phi})\), \( \bar{\lambda}(J^*) \leq \bar{\lambda}(J^*) \) provided

\[
\hat{K}_s \geq K_s \quad \text{and} \quad \frac{\partial \hat{\Phi}_s}{\partial x_s} \leq \frac{\partial \Phi_s}{\partial x_s} \quad \text{for all } s \in S
\]

Theorem 3.5 motivates the following definition of responsiveness. Given two MP-TCP \( \hat{M} \) and \( \hat{M} \), we say that \( \hat{M} \) is more responsive than \( M \) if \( \bar{\lambda}(J^*) \leq \bar{\lambda}(J^*) \). Theorem 3.5(2) implies that an MP-TCP algorithm with a larger \( K_s(x^*_s) \) or more negative definite \( \frac{\partial \Phi_s}{\partial x_s} (x^*_s) \) is more responsive, in the sense that the real parts of the eigenvalues of the Jacobian \( J^* \) have a smaller more negative upper bound.

Then the next result suggests an inevitable tradeoff between responsiveness and friendliness.

Theorem 3.6 (tradeoff): Consider two MP-TCP algorithms \((K, \hat{\Phi})\) and \((\hat{K}, \hat{\Phi})\) with the same gain \( K \). Suppose both satisfy C1–C3 and C5. Then for all \( s \in S \)

\[
\frac{\partial \hat{\Phi}_s(x_s)}{\partial x_s} \leq \frac{\partial \Phi_s(x_s)}{\partial x_s} \quad \Rightarrow \quad \hat{\Phi}_s(x_s) \geq \hat{\Phi}_s(x_s)
\]

In light of Theorems 3.4 and 3.5, Theorem 3.6 says that a more responsive MP-TCP design is inevitably less friendly if they have the same \( K \).
The theorem is easier to understand in the case of SP-TCP, i.e., when \(|s| = 1\) for all \(s \in S\) and \(\Phi_s(x_s) = U'_s(x_s)\). Then it implies that a more concave utility function \(U'_s(x_s)\) has a larger marginal utility, and hence less friendly.

F. Window oscillation

Window oscillations are inherent in loss-based additive increase multiplicative decrease (AIMD) TCP algorithms. We close this section by discussing informally why a larger design \(K_s(x_s)\) generally creates more severe window oscillations. This implies a tradeoff between responsiveness (which is enhanced by a large \(K_s(x_s)\)) and oscillation (which is reduced with a small \(K_s(x_s)\)).

The effect of \(K_s(x_s)\) on window fluctuations can be understood by studying how it affects the decrease \(D_r(w_s)\) per packet loss in the following packet level model:

- For each ACK on route \(r\), \(w_r \leftarrow w_r + I_r(w_s)\).
- For each loss on route \(r\), \(w_r \leftarrow w_r - D_r(w_s)\).

Let \(Z_r \in \{0, 1\}\) be an indicator variable of whether a packet loss is observed on route \(r\) at an arbitrary time in steady state.

\[
D_s(x_s) := \frac{1}{\|x_s\|_1} \mathbb{E} \left( \frac{\sum_{r \in s} D_r(w_s)}{\tau_r} Z_r \sum_{k \in s} Z_k \geq 1 \right)
\]

represents the expected relative reduction in aggregate throughput \(\sum_{r \in s} D_r(w_s)/\tau_r\), given that there is at least one packet loss on some route \(r \in s\). It is a measure of throughput fluctuation for each packet loss that an application experiences. For TCP-NewReno (for which \(s = \{r\}\) and \(w_s\) is a scalar), the window size is halved on each packet loss, \(D_r(w_s) = w_r/2\), and hence \(D_s(x_s) = 1/2\).

To understand \(D_s(x_s)\) for MP-TCP algorithms, we need the following result.

**Lemma 3.1:** Let \(A_1 := \{a_{i1}, a_{i2}, \ldots\}\) with \(|A_1|\) elements. Each element \(a_{ij}\) is an independent binary random variable with \(P(a_{ij} = 1) = 1 - P(a_{ij} = 0) = q_i\). Define \(D_i(A_1) := d_i1(\sum_j a_{ij} \geq 1)\). Then

\[
\mathbb{E} \left( \sum_k D_k(A_k) \right) \sum_{i,j} a_{ij} \geq 1 = \frac{\sum_k d_k q_k |A_k|}{\sum_k q_k |A_k|} + o\left(\sum_k q_k\right)
\]

Suppose each route has a fixed loss probability \(q_r\). Then within each RTT, Lemma 3.1 implies

\[
D_s(x_s) = \frac{1}{\|x_s\|_1} \left( \frac{\sum_{r \in s} w_r q_r D_r(w_s)/\tau_r}{\sum_{r \in s} q_r w_r} + o\left(\sum_{r \in s} q_r\right) \right)
\]

Substituting \(w_r = x_r \tau_r\) and \(x_r D_r(w_s) = \tau_r k_r(x_s)\) from (5), we get, ignoring the high-order terms,

\[
D_s(x_s) = \frac{1}{\|x_s\|_1} \left( \frac{\sum_{r \in s} \tau_r q_r k_r(x_s)}{\sum_{r \in s} \tau_r q_r x_r} \right)
\]

(11)

to the first order. Note that \(k_r(x_s)\) does not affect the equilibrium rates \(x_s\). Hence, with the assumption that \(\tau_r\) are constants, \(D_s(x_s)\) is determined by the functions \(k_r(x_s)\) in steady state.

Specifically an MP-TCP algorithm with a larger \(K_r(x_s)\) tends to have a larger \(D_r(x_s)\) and hence more severe window oscillations. Theorem 3.3 however suggests that a larger \(K_s(x_s)\) also leads to better responsiveness, suggesting an inevitable tradeoff between responsiveness and window oscillation.

IV. IMPLICATIONS AND A NEW ALGORITHM

In this section we discuss the implications of these structural properties on the behavior of existing MP-TCP algorithms. They are further illustrated in experiment results in Section V. The discussion motivates a new design that generalizes the existing MP-TCP algorithm.

A. Implications on existing algorithms

Recall Table I that summarizes the conditions satisfied by the various algorithms. Only EWTCP and Coupled algorithm satisfy C0. Their equilibrium properties can be studied in the standard utility maximization model as done for single-path TCP. Semicoupled and Max algorithms do not satisfy C0 and therefore analysis through utility maximization is not applicable. However Theorem 2.1 below implies that, both Semicoupled and Max algorithms satisfy C1–C3 provided they enable no more than 8 routes. Theorem 3.2 and 3.3 then imply that they have a unique and globally stable equilibrium. It is also easy to show that EWTCP satisfies C1–C3. The Coupled algorithm does not satisfy C2 and is found to have multiple equilibria in [5].

Next we discuss friendliness of existing MP-TCP algorithms. It can be shown that the \(\phi_r(x_s)\) corresponding to these algorithms satisfy:

\[
\phi_{ewtcp}(x_s) \geq \phi_{semicoupled}(x_s) \geq \phi_{max}(x_s) \geq \phi_{coupled}(x_s)
\]

for all \(x_s \geq 0\) if all routes \(r \in s\) have the same round trip time. Since all of them satisfy C4, Theorem 5.4 implies that their friendliness will be in the same order, i.e., their throughputs in the test network of Fig. 1 are ordered as follows:

\[
\text{EWTCP} \geq 1 \geq \text{Semicoupled} \geq \text{Max} \geq \text{Coupled}
\]

This is confirmed by the Linux-based experiment.

Third we will discuss responsiveness of existing MP-TCP algorithms. These algorithms have the same gain function \(k_r(x_s) = 0.5 x_r^2\) and

\[
\frac{\partial \Phi_{ewtcp}}{\partial x_s} \leq \frac{\partial \Phi_{semicoupled}}{\partial x_s} \leq \frac{\partial \Phi_{max}}{\partial x_s} \leq \frac{\partial \Phi_{coupled}}{\partial x_s}
\]

Theorem 3.5 then implies that their responsiveness should be in the same order, as confirmed by our experiments in section V.

Finally we discuss window oscillation of existing MP-TCP algorithms using \(D_s(x_s)\) as the metric. As mentioned in Section III-F, \(D_s(x_s) = 0.5\) for TCP NewReno, a benchmark.

1When \(a < 1\), the MP-TCP source can obtain even smaller throughput than the competing single-path TCP source.
single-path TCP algorithm. According to (11), if $k_r(x_s) \leq 0.5x_r\|x_s\|_1$, we have, to the first order

$$D_s(x_s) \leq \frac{1}{2} \sum_{r\in S} \tau_r q_r x_r \|x_s\|_1 = \frac{1}{2}$$

All existing MP-TCP algorithms have the same $k_r(x_s) = 0.5x_r^2 \leq 0.5x_r\|x_s\|_1$, with strict inequality if $|s| > 1$ and $x_r > 0$ for at least two $r \in s$. Thus enabling MP-TCP always tend to reduce window oscillation for existing algorithms compared to TCP NewReno. Moreover, the window oscillation is always reduced compared to TCP NewReno when $k_r(x_s) \leq 0.5x_r\|x_s\|_1$.

**B. A generalized algorithm**

Consider the class of algorithms parametrized by $(\beta, n, \eta)$ as follows:

$$\begin{align*}
    k_r(x_s) &= \frac{1}{2} x_r (x_r + \eta (\|x_s\|_\infty - x_r)), & \eta \geq 0 \\
    \phi_r(x_s) &= \frac{1}{2} \frac{(1-\beta)x_r + \beta \|x_s\|_1}{\|x_s\|_1^2}, & n \in \mathbb{N}_+, \beta \geq 0 \quad (12)
\end{align*}$$

This class includes the Max ($\beta = 1, \eta = 0, n = \infty$), Coupled ($\beta = 0, \eta = 0$), and Semicoupled ($\beta = 1, \eta = 0, n = 1$) algorithms as special cases when all RTTs on different paths of the same source are the same, i.e., $\tau_r = \tau_s$, $r \in s$.

The next result characterizes a subclass that has a unique and locally stable equilibrium point.

**Theorem 4.1:** Fix any $\eta \geq 0$ and $n \in \mathbb{N}_+$. For any $s \in S$, the $\phi_r(x_s)$ in (12) satisfies

1. C1 if $\beta > 0$.
2. C2-C3 if $0 < \beta \leq 1$, $|s| \leq 8$ and $\tau_r$ are the same for all $r \in s$ (assuming $H$ has full row rank). The requirement that $|s| \leq 8$ is not restrictive since in practice a device may typically enable no more than 3 paths. The requirement that $\tau_r$ are the same for all $r \in s$ is used in proving the negative definiteness of the (symmetric part of the) Jacobian of $\Phi_s(x_s)$. Since a negative definite matrix remains negative definite after small enough perturbations of its entries, Theorem 4.1 holds if the RTTs of the subpaths do not differ much. This (sufficient) condition seems reasonable as two paths between the same source-destination pair often have similar RTTs if both are wireline paths. Note that our experiments in section V show that the RTT of WiFi is usually much smaller than that of 3G.

For the class of algorithms specified by (12), Theorem 4.1 motivates a design space defined by $\beta \in (0, 1], \eta \geq 0, n \in \mathbb{N}_+$, where the control the tradeoff between responsiveness and responsiveness and $\eta$ controls the tradeoff between responsiveness and window oscillation. In Table II we summarize how the parameters $(\beta, \eta, n)$ affect the performance.

We now describe our design philosophy. As discussed above the design of MP-TCP algorithms involves inevitable tradeoffs among responsiveness, friendliness, and the severity of window oscillation. Specifically a design is more responsive if it has a higher gain $K'$ or a more negative definite Jacobian $[\partial\Phi_s/\partial x_s]^{\top}$ (Theorem 3.5). However a larger $K_s$ usually creates a bigger window oscillation; a more negative definite $[\partial\Phi_s/\partial x_s]^{\top}$ implies a larger $\Phi_s$, usually hurting friendliness

| Performance | Parameter | Parameters in (12) |
|-------------|-----------|--------------------|
| TCP friendliness | $\phi_s(x_s)$ | $\beta \uparrow, n \uparrow$ |
| Responsiveness | $k_r(x_s)$ | $\beta \uparrow, n \uparrow, \eta \uparrow$ |
| Window oscillation | $k_r(x_s)$ | $\eta \downarrow$ |

**TABLE II:** How design choices affect MP-TCP performance.

Fig. 2: Network for our Linux-based experiments on TCP friendliness and responsiveness, with $N_1$ single-path TCP flows and $N_2$ single-path TCP flows sharing 2 links of capacity $c_1, c_2$ and delay $\tau_1, \tau_2$. MP-TCP flows maintain two routes with rate $x_1, x_2$. Single-path TCP flows maintain one route with rate $x_3$. (Theorems 3.6 and 3.4). This is summarized in Table II. Since enabling multiple paths already reduces window oscillation compared to single-path TCP (section IV-A), MP-TCP can afford to use a relatively large gain $K_s$ for responsiveness. This does not compromise too much on window oscillation, but allows us to use a less negative definite Jacobian $[\partial\Phi_s/\partial x_s]^{\top}$ with a smaller $\Phi_s$ to maintain sufficient TCP friendliness. Moreover, responsiveness is mainly affected by subpaths with small throughput while window oscillation is mainly affected by subpaths with large throughput. The parameter $\eta$ in the generalized algorithm (12) scales $k_r(x_s)$ in the right way: a path $r$ that has a large $x_r$ has $k_r(x_s) \approx 0.5x_r^2$ and hence a similar degree of window oscillation as existing algorithms, while a path $r$ with a small $x_r$ has larger $k_r(x_s)$ than that under a design with zero $\eta$ and therefore is more responsive.

Our experiments show that Max algorithm $(\beta, \eta, n) = (1, 0, \infty))$ overtakes too much of the competing single-path TCP flows. Hence, we can only use a smaller $\beta$ since $n$ is already infinite in order to improve friendliness. To compensate the responsiveness performance, we will use a larger $\eta$, which will sacrifice window oscillation performance.

The algorithm given at the end of Section II corresponds to the choice $(\beta, \eta, n) = (0.2, 0.5, \infty)$. Instead of allowing the window size to drop to 1 for a packet loss, we add a cap for the decrement of window size, which improves the performance according to our experiments. Note that there is no "best" parameter settings since there are tradeoffs among all the performance metrics and we choose $(\beta, \eta, n) = (0.2, 0.5, \infty)$ based on our experiments in section V which show that this parameter setting strikes a good balance among responsiveness, friendliness, and window oscillation.

**V. EXPERIMENT**

In this section we summarize our experimental results that illustrate the above analysis. In addition to the MP-TCP algorithms illustrated in section II-B we also include
the recently developed OLIA MP-TCP algorithm [8]. We evaluate the MP-TCP algorithms using a reference Linux implementation of MP-TCP, Multipath TCP v0.88 [13]. Since it currently includes only Max and OLIA algorithms, we implement EWTCP, Semicoupled, Coupled, and our algorithm in the reference implementation. For the Coupled and our algorithm, the minimum *sthresh* is set to 1 instead of 2 when more than 1 path is available.

The network topology is shown in Fig. 2. In the testbed, all nodes are Linux machines with a quad-core Intel i5 3.33GHz processor, 4GB RAM and multiple 1Gbps Ethernet interfaces, running Ubuntu 13.10 (Linux kernel 3.11.8). The network parameters such as $c_1$, $c_2$, $\tau_1$, and $\tau_2$ are controlled by Dummynet [14].

Our experiments are divided into three parts. First we compare TCP friendliness of our algorithm and prior algorithms. The result confirms that the Coupled algorithm is the friendliest, OLIA and ours are close to the Coupled algorithm and friendlier than the other algorithms. Second we compare the responsiveness of each algorithm in a dynamic environment where flows come and go. The result shows that the Coupled and OLIA algorithms are unresponsive (illustrating the tradeoff between responsiveness and friendliness). EWTCP is the most responsive; ours is similar in responsiveness but friendlier than single-path TCP. Finally we show that all MP-TCP algorithms have smaller average window oscillations than single-path TCP.

These experiments confirm our analytical results and suggest our design choice strike a good balance among friendliness, responsiveness, and window oscillation.

A. TCP friendliness

We study TCP friendliness of each algorithm, first with paths of similar RTTs and then with paths of different RTTs, which emulates the wireless scenario. We assume all the flows are long lived and focus on the steady state throughput.

In the first set of experiments, we let $\tau_1 = \tau_2 = 5$ms, $c_1 = c_2 = 10$Mbps and $N_1 = N_2 = 5$. The average aggregate throughput of MP-TCP and single-path TCP users are shown in Table III. The EWTCP, Semicoupled and Max algorithms are significantly less friendly to SP-TCP. The Coupled algorithm is the friendliest and OLIA and our algorithm are very close.

|       | ewtcp | semi. | max | ours | coupled | olia |
|-------|-------|-------|-----|------|---------|------|
| MP-TCP | 2.858 | 2.66  | 2.58 | 2.472 | 2.36    | 2.488 |
| SP-TCP | 0.875 | 1.09  | 1.16 | 1.274 | 1.39    | 1.26  |

In the second set of experiments, we assume a highly heterogeneous RTTs by emulating the scenario of a mobile device with both 3G and WiFi access. WiFi access usually has higher capacity and lower delay compared to 3G. Specifically, we set $\tau_1 = 10$ms, $c_1 = 8$Mbps for the first link to emulate WiFi access and $\tau_2 = 100$ms, $c_2 = 2$Mbps for the second link to emulate 3G access. When there exists single-path TCP flows, i.e., $N_2 > 0$, the behaviors of all the algorithms are similar to the equal RTT case in the first set of simulation. The EWTCP, Semicoupled and Max algorithms are significantly less friendly to SP-TCP. The Coupled algorithm is the friendliest and OLIA and our algorithm are very close. However, when there is no single-path TCP flow, i.e., $N_1 = 1$ and $N_2 = 0$, OLIA is not effective to take all the available capacity while the other algorithms do not have such problem. The average throughput of MP-TCP user is shown in Table IV.

|       | ewtcp | semi. | max | ours | coupled | olia |
|-------|-------|-------|-----|------|---------|------|
| MP-TCP | 9.24  | 9.22  | 9.24 | 9.28 | 9.23    | 8.94 |

B. Responsiveness

We use the network in Fig. 2 with $c_1 = c_2 = 20$Mbps, $\tau_1 = \tau_2 = 10$ms and $N_1 = N_2 = 5$. To demonstrate the dynamic performance of each algorithm, we assume the MP-TCP flow is long lived while the SP-TCP flows start at 40s and end at 80s. We record the aggregate throughput of the SP-TCP flows from 40-80s, which measures the friendliness of MP-TCP. We also measure the time for the congestion window on the second path to recover of MP-TCP users. It measures the responsiveness of MP-TCP. These measurements are shown in Table V and the congestion window and throughput trajectories of all algorithms are shown in Fig. 4.

|       | ewtcp | semi. | max | ours | coupled | olia |
|-------|-------|-------|-----|------|---------|------|
| Convergence | 3.25 | 7.46 | 17.75 | 14.73 | 94.36 | 58.5 |
| SP-TCP | 13.89 | 15.35 | 15.8 | 16.28 | 16.64 | 16.97 |

EWTCP is the most responsive among all the algorithms. Ours is as responsive as the Max algorithm, yet significantly friendlier than EWTCP. Both Coupled and OLIA algorithms take an excessively long time to recover. For Coupled algorithm, the excessively slow recovery of the congestion window on the second path (see Fig. 4) is due to the design that increases the window roughly by $w_r/(\sum_{k \in S} w_k)^2$ on each ACK assuming the RTTs are similar. After the single-path TCP flow has left, $w_2$ is small while $w_1$ is large, so that $w_2/(w_1 + w_2)^2$ is very small. It therefore takes a long time for $w_2$ to increase to its steady state value. In general, under the Coupled algorithm, a route with a large throughput can greatly suppress the throughput on another route even though the other route is underutilized. The reason of the poor responsiveness performance of OLIA can be explained using similar argument as Coupled algorithm since they have the same increment/decrement for each ACK/loss in this scenario.
C. Window oscillation

We use a single-link network model to compare window oscillation under MP-TCP and single-path TCP. First a MP-TCP flow initiates two subpaths through that link, and we measure the window size of each subpath and their aggregate window size. Then a TCP-Reno flow traverse the same link and we measure its window size. The results are shown in Fig. 3 for our algorithm in comparison with single-path TCP (other MP-TCP algorithms have a similar behavior). They confirm that enabling multiple paths reduces the average window oscillation compared with only using single path.

![Fig. 3: Window oscillation: the red trajectories represent throughput fluctuations experienced by the application in the case of MP-TCP and the case of single-path TCP.](image)

VI. CONCLUSION

We have presented a model for MP-TCP and identified designs that guarantee the existence, uniqueness and stability of the network equilibrium. We have characterized the design space and study the tradeoff among TCP friendliness, responsiveness, and window oscillation. We have proposed a new MP-TCP algorithm that generalizes existing algorithms and strikes a good balance among these properties. We have implemented our algorithm in the Linux kernel and used it to evaluate the performance of our algorithm.

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![Fig. 4: Responsiveness Performance: congestion window trajectory of MP-TCP for each path (left column). SP-TCP starts at time 40s and ends at 80s. The throughput of SP-TCP and total throughput of MP-TCP are shown in the right column. Parameters: $\tau_1 = \tau_2 = 10$ms, $c_1 = c_2 = 20$Mbps and $N_1 = 1, N_2 = 5$.](image)
where the algorithm defined by (3)–(4) has an underlying utility function

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APPENDIX A

PROOF OF THEOREM 3.1 (UTILITY MAXIMIZATION)

The Lagrangian of (9) is:

\[ L(x, p) = \sum_{s \in S} U_s(x_s) - \sum_{l \in L} p_l(y_l - c_l) \]

\[ = \sum_{s \in S} U_s(x_s) - \sum_{l \in L} p_l(\sum_{r \in R} H_{lr}x_r - c_l) \]

\[ = \sum_{s \in S} \left( U_s(x_s) - \sum_{r \in s} x_rq_r \right) + \sum_{l \in L} p_lc_l \]

where \( p \geq 0 \) are the dual variables and \( q_r := \sum_{r \in R} H_{lr}p_l \).

Then the dual problem is

\[ D(p) = \sum_{s \in S} \max_{x_r \geq 0} \left\{ B_s(x_s, p) \right\} + \sum_{l \in L} p_lc_l \quad p \geq 0 \]

where \( B_s(x_s, p) = U_s(x_s) - \sum_{r \in s} x_rq_r \). The KKT condition implies that, at optimality, we have

\[ \frac{\partial U_s(x_s)}{\partial x_r} < q_r \Rightarrow x_r = 0 \quad \text{and} \quad x_r > 0 \Rightarrow \frac{\partial U_s(x_s)}{\partial x_r} = q_r \]

\[ y_l < c_l \Rightarrow p_l = 0 \quad \text{and} \quad p_l > 0 \Rightarrow y_l = c_l \]

Comparing with (6)–(7) we conclude that, if a MP-TCP algorithm defined by (3)–(4) has an underlying utility function \( U_s \), then we must have

\[ \frac{\partial U_s(x_s)}{\partial x_r} = \phi_r(x_s) \quad r \in s, \ x_r > 0 \]

Given \( \phi_r(x_s) \), (15) has a continuously differentiable solutions \( U_s(x_s) \) if and only if the Jacobian of \( \Phi_s(x_s) \) is symmetric, i.e., if and only if

\[ \frac{\partial \Phi_s(x_s)}{\partial x_s} = \left[ \frac{\partial \Phi(x_s)}{\partial x_s} \right]^T \]

APPENDIX B

PROOF OF THEOREM 3.2 (EXISTENCE AND UNIQUENESS)

A. Proof of part 1

For any link \( l \in L \), let

\[ p_{-l} = \left\{ p_1, \ldots, p_{l-1}, p_{l+1}, \ldots, p_{|L|} \right\}, \]

whose component composes of all the elements in \( p \) except \( p_l \).

\[ g_l(p) := c_l - \sum_{r \in R} x_r = c_l - \sum_{r : r \in s, \ r \in R} y^*_l(p_l, p_{-l}) \]

and \( h_l(p) := -g^2_l(p) \).

According to C1, we have the following two facts, which will be used in the proof.

- \( g_l(p) \) is a nondecreasing function of \( p_l \) on \( \mathbb{R}_+ \) since \( y^*_l(p) \) is a nonincreasing function of \( p_l \).
- \( \lim_{p_l \to \infty} g_l(p_l, p_{-l}) = c_l \) since \( y^*_l(p) \) is a nonincreasing function of \( p_l \).

Next, we will show that \( h_l(p) \) is a quasi-concave function of \( p_l \). In other words, for any fixed \( p_{-l} \), the set \( S_a := \{ p_l \mid h_l(p) \geq a \} \) is a convex set. If \( g_l(0, p_{-l}) \geq 0 \), then

\[ g_l(p_l, p_{-l}) \geq g_l(0, p_{-l}) \geq 0 \forall p_l \geq 0, \]

which means \( h_l(p_l, p_{-l}) \) is a nonincreasing function of \( p_l \), hence is a quasi-concave function of \( p_l \) and

\[ \arg \max_{p_l} h_l(p_l, p_{-l}) = 0. \] (16)

On the other hand, if \( g_l(0, p_{-l}) < 0 \), then there exists a \( p^*_l > 0 \) such that \( g_l(p^*_l, p_{-l}) = 0 \) since \( y^*_l(p) \) is continuous and \( \lim_{p_l \to \infty} g_l(p_l, p_{-l}) = c_l > 0 \). Note that \( g_l(p) \) is a non-decreasing function of \( p_l \), then \( h_l(p_l, p_{-l}) \) is nondecreasing for \( p_l \in [0, p^*_l] \) and nonincreasing for \( p_l \in [p^*_l, \infty) \). Hence, \( h_l(p_l, p_{-l}) \) is also a quasi-concave function of \( p_l \) in this case and

\[ \max_{p_l} h_l(p_l, p_{-l}) = 0. \] (17)

By Nash theorem, if \( h_l(p_l, p_{-l}) \) is a quasi-concave function of \( p_l \) for all \( l \in L \) and \( p \) in a bounded set, then there exists a \( p^* \in \mathbb{R}_+^{|L|} \) such that

\[ p^*_l = \arg \max_{p_l \in \mathbb{R}_+} h_l(p_l, p^*_{-l}). \]

According to (16) and (17), for any \( l \in L \), either \( p^*_l > 0 \) or \( g^*_l(p^*) > 0 \) but not both holds at any time. Therefore \( p^* \) satisfies Eqn. (7). Since \( q = R^T p \), there exists an \( x^* \) to (6). Hence there exists at least one solution \( (x, p) \) that satisfies (6) and (7).

B. Proof of part 2

Lemma B.1: Assume a function \( F : \mathbb{R}^n \to \mathbb{R}^n \) is continuously differentiable and \( \left[ \frac{\partial F}{\partial x} (x) \right]^T \) is negative definite for all \( x \). Then for any \( x_1 \neq x_2 \in \mathbb{R}^n \),

\[ (x_1 - x_2)^T (F(x_1) - F(x_2)) < 0. \]

Proof: Fix any \( x_1 \neq x_2 \in \mathbb{R}^n \). Define \( A(t) := F (tx_1 + (1-t)x_2) \). Since \( \partial F / \partial x \) is continuous, there exists
a \lambda < 0 such that the eigenvalues of $[\partial F/\partial x]^T \lambda$ over the compact set \{$tx_1 + (1-t)x_2 | 0 \leq t \leq 1$\}. Then

$$
(x_1 - x_2)^T (F(x_1) - F(x_2)) = \int_0^1 (x_1 - x_2)^T \frac{dA}{dt} (\tau) d\tau
$$

$$
\leq \lambda \|x_1 - x_2\|^2 \leq 0
$$

**Lemma B.2:** Suppose C3 holds. Then $x^*_r > 0$ at equilibrium for all $r \in R$.

**Proof:** Suppose $x^*_r = 0$. Then $q^*_r \geq \phi_r(x^*_r) = \infty$ by C3 and hence there is a link $l \in R$ with $p^*_l = \infty$. But then, for all paths $r' \ni l$, $q^*_{r'} = \infty$ and hence $x^*_{r'} = 0$ by C3. This implies $y^*_l = 0 < c_l$, and hence $p^*_l = 0$ by \(\tau\), contradicting $p^*_l = \infty$.

Recall the vector notations that $x := (x_s, s \in S) := (x_r, r \ni s, s \in S)$ and $\Phi(x) := (\Phi_s(x_s), s \in S) := (\Phi_r(x_s), r \ni s, s \in S)$. To prove uniqueness of the equilibrium, suppose for the sake of contradiction that there are two distinct equilibrium points $(x, p)$ and $(\hat{x}, \hat{p})$. By Lemma B.2, we have $x > 0$ and $\hat{x} > 0$. Hence \(\Phi(x) = q = H^T p\) and $\Phi(\hat{x}) = \hat{q} = H^T \hat{p}$. By Lemma B.1 and assumption C2, we then have

$$
0 > (x - \hat{x})^T (\Phi(x) - \Phi(\hat{x}))
$$

$$
= (x - \hat{x})^T H^T (p - \hat{p})
$$

$$
= (p - \hat{p})^T (y - \hat{y})
$$

Hence

$$
p^T y + \hat{p}^T \hat{y} < p^T \hat{y} + \hat{p}^T y
$$

Equilibrium condition \(\tau\) implies

$$
p^T (c - y) = 0 \quad \text{and} \quad \hat{p}^T (c - \hat{y}) = 0
$$

$$
y \leq c \quad \text{and} \quad \hat{y} \leq c
$$

Substituting \(19\) into \(18\) yields

$$
p^T(c - y) + \hat{p}^T(c - \hat{y}) < p^T \hat{y} + \hat{p}^T y
$$

$$
p^T(c - \hat{y}) + \hat{p}^T(c - y) < 0
$$

But \(20\) implies that the left-hand side of the last inequality is nonnegative (since $p \geq 0$, $\hat{p} \geq 0$), a contradiction. Hence the equilibrium is unique.

**APPENDIX C**

**PROOF OF THEOREM 3.3 (STABILITY)**

We will construct a Lyapunov function and use LaSalle’s invariance principle \(15\) to prove global asymptotic stability of the unique equilibrium point $(x^*, p^*)$. Define $\delta x := x - x^*$, $\delta p := p - p^*$. Consider the candidate Lyapunov function:

$$
V(x, p) = \sum_{r \in R} \int_{x_r^*}^{x_r} \frac{\mu}{k_r(z)} dz + \frac{1}{2} \sum_{l \in L} \delta p_l^2
$$

By definition, $V(x, p) > 0$ for all $(x, p) \neq (x^*, p^*)$ and $V(x, p) = 0$ if $(x, p) = (x^*, p^*)$. Furthermore $V$ is radially unbounded, i.e., $V(x, p) \rightarrow \infty$ as $\|x, p\| \rightarrow \infty$. Finally

$$
\dot{V}(x, p) = \sum_{r \in R} \frac{1}{k_r(x_r)} \delta x_r \dot{x}_r + \sum_{l \in L} \frac{1}{\gamma_l} \delta p_l \dot{p}_l
$$

If $\delta x_r \neq 0$ then we have (since $k_r(x_s) = k_r(x_r)$)

$$
\frac{1}{k_r(x_r)} \delta x_r \dot{x}_r = \delta x_r (\phi_r(x_s) - q_r^+)^-
$$

$$
\leq \delta x_r (\phi_r(x_s) - \phi_r(x_s) - q_r^+)
$$

The first inequality holds since $(\phi_r(x_s) - q_r^+) = \phi_r(x_s) - q_r$ if $x_r > 0$ and $\phi_r(x_s) - q_r \leq 0$, $\delta x_r = -x^*_r$ if $x_r = 0$. The last equality holds since $\phi_r(x^*_s) = q^*_r$ by Lemma B.2 and \(\tau\).

$$
\sum_{r \in R} \frac{1}{k_r(x_r)} \delta x_r \dot{x}_r \leq -\delta x^T (\Phi(x) - \Phi(x^*)) - \delta x^T \delta q
$$

$$
< -\delta x^T H^T \delta \Phi
$$

where the last inequality holds since $\delta x^T (\Phi(x) - \Phi(x^*)) < 0$ by Lemma B.1 and assumption C2. Similarly

$$
\frac{1}{\gamma_l} \delta p_l \dot{p}_l = \delta p_l (y_l - c_l^+)^- \leq \delta p_l (y_l - c_l) \leq \delta p_l \delta y_l
$$

where the last inequality holds since $\delta p_l c_l \geq \delta p_l y_l^*$ by the equilibrium condition \(\tau\).

$$
\sum_{l \in L} \frac{1}{\gamma_l} \delta p_l \dot{p}_l \leq \delta p^T H \delta x
$$

Therefore if $\delta x \neq 0$ then

$$
\dot{V}(x, p) < -\delta x^T H^T \delta p + \delta p^T H \delta x = 0
$$

and if $\delta x = 0$ then $\dot{V}(x, p) = 0$. This means $\dot{V}(x, p) \leq 0$ and $V$ is indeed a Lyapunov function.

Consider the set

$$
Z := \{ (x(t), p(t)) | V(x(t), p(t)) = 0 \text{ for all } t \geq 0 \}
$$

of trajectories on which $\dot{V} = 0$. If the only trajectory in $Z$ is the trivial trajectory $(x, p) = (x^*, p^*)$ then LaSalle’s invariance principle implies that $(x^*, p^*)$ is globally asymptotically stable. We now show that this is indeed the case.

As shown above $\dot{V} \equiv 0$ implies $\delta x \equiv 0$, i.e., any trajectory $(x(t), p(t))$ in $Z$ must have $x(t) = x^*$ for all $t \geq 0$. This means $x \equiv 0$ and hence, for all $t \geq 0$, $q(t) = \Phi(x(t))$ since $x(t) = x^* > 0$ by Lemma B.2. That is, for all $t \geq 0$, $H^T p(t) = \Phi(x^*)$ and hence $p(t) = p^*$ since $H$ has full row rank by C3. Therefore $(x, p) \equiv (x^*, p^*)$ is indeed the only trajectory in $Z$. This completes the proof of Theorem 3.3.

**APPENDIX D**

**PROOF OF THEOREM 3.4 (FRIENDLINESS)**

Let the MP-TCP source be defined by

$$
\phi_r(x_s; \mu) = \mu \phi_r(x_s) + (1 - \mu) \phi_r(x_s), \quad \mu \in [0, 1]
$$

Algorithm $\tilde{M}$ and $\tilde{M}$ corresponds to $\mu = 0$ and $\mu = 1$ respectively. Let $x_g$ and $r_g$ be the throughput and RTT of the
TCP NewReno source in Fig. 1. The equilibrium is defined by $F(x, \mu) = 0$ where $x := (x_s, x_g)$ and $F$ is given by:

$$
\Phi_s(x_s; \mu) - \frac{1}{\gamma^2 x_s^2} = 0
$$

$$
1^T x_s + x_g = c
$$

where the first equation follows from

$$
p^* = \frac{1}{\gamma^2 x_s^2} = \phi_s(x_s; \mu), \quad r \in s
$$

and $p^*$ is the congestion price at the bottleneck link. Applying the implicit function theorem, we get

$$
\frac{dx}{d\mu} = -\left(\frac{\partial F}{\partial x}\right)^{-1} \frac{\partial F}{\partial \mu}
$$

where the inverse exists by condition C2. C2 also guarantees the inverse of $\frac{\partial \Phi_s}{\partial x_s}(x_s; \mu)$, denoted by $D(\mu)$; C4 ensures $\sum_{i \in s} D_{i,j}(\mu) \leq 0$. Let

$$
A := \frac{\partial \Phi_s}{\partial x_s} - \frac{2}{x_g^2} 11^T \quad \text{and} \quad d := 1 - \frac{2}{x_g^2} \sum_{i,j} D_{i,j}(\mu)
$$

Then

$$
\begin{bmatrix}
\frac{\partial \Phi_s}{\partial x_s} & 0 & \frac{2}{x_g^2} 1
\end{bmatrix}^{-1} = \begin{bmatrix}
A^{-1} & -D1d \\
-d1^TA^{-1} & d^{-1}
\end{bmatrix}
$$

Thus

$$
1^T \frac{\partial x_s}{\partial \mu} = -1^T 0 \left(\frac{\partial F}{\partial x}\right)^{-1} \frac{\partial F}{\partial \mu}
$$

By matrix inverse formula,

$$
A^{-1} = \left(\frac{\partial \Phi_s}{\partial x_s} - \frac{2}{x_g^2} 11^T\right)^{-1}
$$

$$
= D(\mu) + \frac{1}{x_g^2 - 1^T D(\mu) 1} D(\mu) 11^T D(\mu)
$$

Substitute it into (21), we have

$$
1^T A^{-1}(\Phi_s(x_s) - \Phi_s(x_s)) = 1 + \frac{1}{x_g^2 - 21^T D(\mu) 1} \sum_{i \in s} D_{i,r}(\mu) \frac{\partial \Phi_s}{\partial x_r}(x_s) - \frac{\partial \Phi_s}{\partial x_r}(x_s)
$$

where the inequality follows because $D(\mu)$ is negative definite, $\sum_{i \in s} D_{i,r}(\mu) < 0$ and $\frac{\partial \Phi_s}{\partial x_r}(x_s) - \frac{\partial \Phi_s}{\partial x_r}(x_s) \geq 0$. Thus we have

$$
1^T A^{-1}(\Phi_s(x_s) - \Phi_s(x_s)) \leq 0
$$

**Appendix E**

**Proof of Theorem 3.5 (Responsiveness)**

**A. Proof of part 1**

Fix any eigenvalue $\lambda$ of $J^*$. Let $z := (x, p) \in Z$ be the corresponding eigenvector with $\|z\|_2 = 1$. Then we have

$$
\lambda \begin{bmatrix}
\tilde{\Lambda}^{-1} \\
\tilde{\Omega}_r
\end{bmatrix}
= \begin{bmatrix}
\frac{\partial \Phi_s}{\partial x} & -H^T \\
H & 0
\end{bmatrix}
\begin{bmatrix}
x \\
p
\end{bmatrix}
$$

Hence

$$
\lambda \begin{bmatrix}
\tilde{\Lambda}^{-1} \\
\tilde{\Omega}_r
\end{bmatrix}
= \begin{bmatrix}
\frac{\partial \Phi_s}{\partial x} & -H^T \\
H & 0
\end{bmatrix}
\begin{bmatrix}
x \\
p
\end{bmatrix}
$$

Premultiplying $z^H$ on both sides, we have

$$
\lambda = \frac{\text{Re}(\lambda) \lambda^+ - \text{Im}(\lambda)^+ x^H H x - x^H H^T p}{x^H \tilde{\Lambda}^{-1} x + p^H \tilde{\Omega}_r}
$$

The denominator is real and positive, and $(p^H H x - x^H H^T p)$ in the numerator is imaginary. Hence

$$
\text{Re}(\lambda) = \frac{\text{Re}(x^H \frac{\partial \Phi_s}{\partial x})}{x^H \tilde{\Lambda}^{-1} x + p^H \tilde{\Omega}_r}
$$

$$
= \frac{x^H \tilde{\Lambda}^{-1} x + p^H \tilde{\Omega}_r}{< 0}
$$

where the last inequality holds because the numerator is negative by condition C2 and the denominator is positive. Since this holds for all eigenvalues $\lambda$ of $J^*$, the linearized system is stable. Moreover $\text{Re}(\lambda) \leq \chi(J^*) \leq 0$ as desired.

**B. Proof of part 2**

Consider two MP-TCP algorithms $(\tilde{K}, \tilde{\Phi})$ and $(\hat{K}, \hat{\Phi})$ such that

$$
\hat{K}_s \geq \tilde{K}_s \quad \text{and} \quad \frac{\partial \hat{\Phi}_s}{\partial x_s} \leq \frac{\partial \tilde{\Phi}_s}{\partial x_s} \quad \text{for all} \quad s \in S
$$

For any (nonzero) $z = (x, p) \in Z$ we have

$$
0 \leq x^H \tilde{\Lambda}_s^{-1} x \leq x^H \hat{\Lambda}_s^{-1} x
$$

$$
x^H \frac{\partial \hat{\Phi}_s}{\partial x} x \leq x^H \hat{\Lambda}_s^{-1} x < 0
$$

Hence $\chi(J^*) \leq \chi(J^*)$.

**Appendix F**

**Proof of Theorem 3.6 (Tradeoff)**

Fix an $s$. Let $f_r(x_s) := \hat{\phi}_r(x_s) - \tilde{\phi}_r(x_s)$ and $F(x_s) := (f_r(x_s), r \in s) = \Phi_s(x_s) - \Phi_s(x_s)$. Suppose for the sake of contradiction that $\frac{\partial \Phi_s(x_s)}{\partial x_s} \leq \frac{\partial \hat{\Phi}_s(x_s)}{\partial x_s}$ but $\hat{\phi}_s(x_s) \geq \tilde{\phi}_s(x_s)$ does not hold, i.e., there exists a finite $x_s^0$ and a $r \in s$ such that

$$
f_r(x_s^0) = \hat{\phi}_r(x_s^0) - \tilde{\phi}_r(x_s^0) < 0
$$

Since $\frac{\partial F}{\partial x^0} \geq 0$ by assumption, a trivial modification of Lemma B.1 shows that, for all $x_s \neq x_s^0$, $(x_s - x_s^0)^T (F(x_s) - F(x_s^0)) \leq 0$, i.e.,

$$
0 \geq \sum_{r \in s} (x_r - x_r^0)(f_r(x_s) - f_r(x_s^0))
$$

(26)
Choose an \( x_s \) as follows: for all \( r' \neq r \), choose \( x_{r'} = x^0_r \), and then use condition C5 to choose an \( x_r < \infty \) large enough so that \( x_r > x^0_r \) and \( f_r(x_s) > f_r(x^0_s)/2 \). With this \( x_s \), \( (26) \) becomes
\[
0 \geq (x_r - x^0_r) \left( f_r(x_s) - f_r(x^0_s) \right) > (x_r - x^0_r) \left( -\frac{f_r(x^0_s)}{2} \right) > 0
\]
where the last inequality follows from \( (25) \). This is a contradiction and hence \( \Phi_s(x_s) \geq \Phi_s(x_s) \).

**Appendix G**

**Proof of Theorem 4.1**

We will show the results hold for any \( n \in \mathbb{N}_+ \). Since \( \lim_{n \to \infty} \| x_s \|_n = \| x_s \|_\infty \), the results also hold for \( n = \infty \). When \( \beta = 0 \), it is easy to show that \( \phi_r \) satisfies C1 and
\[
\begin{bmatrix}
\phi_r \\
\sqrt{C}
\end{bmatrix}^T
\]
is negative semidefinite under the conditions of the theorem. We hence prove the theorem for \( \beta > 0 \).

**A. Proof of part 1**

Fix any \( n \in \mathbb{N}_+ \) and \( \beta > 0 \). Fix any finite \( p \geq 0 \) such that \( q_r > 0 \) for all \( r \). Fix any \( s \in S \). We now show that there exists an \( x_s \geq 0 \) that satisfies \( (24) \), in particular \( \phi_r(x_s) = q_r \), in two steps.

First, there exists an \( x_s \) that satisfies \( \phi_r(x_s) = q_r \) if and only if
\[
\phi_r(x_s) = \frac{2}{\tau^r_s \| x_s \|_1^2} \left( 1 + \beta \left( \frac{\| x_s \|_n}{x_r} - 1 \right) \right) = q_r, \tag{27}
\]
which is equivalent to
\[
x_r = \frac{2\beta}{\| x_s \|_n} = \frac{2\beta + q_r \tau^r_s}{2\beta + q_r \tau^r_s \| x_s \|_1^2 - 2} \tag{28}
\]
Since this holds for all \( r \in s \), we have
\[
1 = \sum_{r \in s} \left( \frac{x_r}{\| x_s \|_n} \right)^n = \sum_{r \in s} \left( \frac{2\beta}{2\beta + q_r \tau^r_s \| x_s \|_1^2 - 2} \right)^n =: \psi (\| x_s \|_1^2)
\]
Clearly \( \psi (C) \to 0 \) as \( C \to \infty \). Let
\[
\tilde{C} := \frac{2}{\min_{r \in s} q_r \tau^r_s} \tag{30}
\]
Then \( C < \infty \) since \( q_r > 0 \) for all \( r \) by assumption. Moreover \( q_r \tau^r_s \tilde{C} \geq 2 \) for all \( r \in s \) and hence
\[
\psi (C) = 1 + \sum_{r \notin s} \left( \frac{2\beta}{2\beta + q_r \tau^r_s \tilde{C} - 2} \right)^n > 1
\]
where \( \tilde{r} \) is a minimizing \( r \in s \) in \( (30) \). Since \( \psi (C) \) is continuous, there exists an \( \tilde{C} \in [C, \infty) \) with \( \psi (\tilde{C}) = 1 \). Moreover such a \( \tilde{C} \) is unique since \( \psi (C) \) is strictly decreasing.

Finally consider the set of \( x_s \) with \( \| x_s \|_1^2 = \tilde{C} \). All such \( x_s \) satisfy \( (28) \) with
\[
x_r = \frac{2\beta}{2\beta + q_r \tau^r_s \tilde{C} - 2} \| x_s \|_n =: a_r \| x_s \|_n \tag{31}
\]
But \( \hat{C} = \| x_s \|_n^2 = \left( \sum_{r \in s} a_r \| x_s \|_n \right)^2 \), implying
\[
\| x_s \|_n = \sqrt{\frac{\tilde{C}}{\sum_{r \in s} a_r}}
\]
In summary, given any finite \( p \geq 0 \) such that \( q_r > 0 \) for all \( r \), a solution \( x_s > 0 \) to \( (28) \) is uniquely given by
\[
x_r = \frac{a_r}{\| x_s \|_n} \sqrt{\tilde{C}}, \quad r \in s \tag{32}
\]
where
\[
a_r := \frac{2\beta}{2\beta + q_r \tau^r_s \tilde{C} - 2}
\]
and \( \hat{C} = \| x_s \|_n^2 \) is the unique value at which \( \psi (\hat{C}) = 1 \).

We now prove the other conditions in C1:
\[
\frac{\partial y^r_t (p)}{\partial p_l} \leq 0, \quad \lim_{p_l \to 0} y^r_t (p) = 0
\]
According to \( (22) \), we can show that \( \hat{C} \) is a decreasing function of \( q_r \) and \( q_r \tau^r_s \hat{C} \) is an increasing function of \( q_r \) for \( r \in s \). Thus, \( \hat{C} \) is a decreasing function of \( p_l \) and \( q_r \tau^r_s \hat{C} \) is an increasing of \( p_l \) if \( l \in r \) because \( q_r = \sum_{l \in L} H_{rl} p_l \). For each \( l \in L \), let \( s_l := \{ r \mid l \in r, r \in s \} \), then by definition and \( (32) \), we have
\[
y^r_t (p) = \sum_{r \in s_l} a_r \sqrt{\hat{C}} = \sum_{r \in s_l} \frac{a_r}{\| x_s \|_n} \sqrt{\tilde{C}}.
\]
Since \( a_r \) is a decreasing function of \( q_r \tau^r_s \hat{C} \), it is also a decreasing function of \( p_l \) if \( l \in r \). Recall that \( \sqrt{\tilde{C}} \) is also a decreasing function of \( p_l \). \( \sqrt{\tilde{C}} \) is thus a decreasing function of \( p_l \), in other words, \( \frac{\partial \psi (p)}{\partial p_l} \leq 0 \).

On the other hand, as \( p_l \to \infty \), \( q_r \to \infty \) for all paths \( r \) traversing \( l \). Then \( x_r \to 0 \) by \( (27) \) for \( l \in r \), which shows \( \lim_{p_l \to \infty} y^r_t (p) = 0 \).

**B. Proof of part 2**

To prove \( \phi_r(x_s) \) satisfies C2 and C3 for \( \beta > 0 \), we will show that the Jacobian \( \partial \Phi_s (x_s) / \partial x_s \) is negative definite if \( 0 < \beta \leq 1 \), \( |s| \leq 8 \) and \( \tau_r \) are the same for all \( r \). Other properties of C2 and C3 are easy to prove and we omit the proof. Fix an \( s \) and let \( \tau_r = \tau \), the common round-trip time for all \( r \in s \).

Let \( A_s := \text{diag} \{ x_s \} \) and
\[
a_s := \left( \frac{2x_r}{\| x_s \|_1} - \frac{x^0_r}{\| x_s \|_n} \right), \quad r \in s \tag{33}
\]
Then the Jacobian of \( \Phi_s \) at \( x_s \) is
\[
\frac{\partial \Phi_s}{\partial x_s} = -\frac{4(1 - \beta)}{\tau^s \| x_s \|_1^2} 1^T + 2 - 2\frac{\| x_s \|_n}{\tau^2 \| x_s \|_1} \Lambda_s^{-1} (I_s + 1a_s^T) \Lambda_s^{-1}
\]
and it is negative definite for \( \beta > 0 \) if \( (I_s + 1a_s^T) \) is positive definite. We now show that this is indeed the case when \( |s| \leq 8 \), i.e., for any \( z_s \in \mathbb{R}^{|s|} \),
\[
z_s^T (I_s + 1a_s^T) z_s = \| z_s \|_2^2 + \sum_{r \in s} z_r \sum_{r \in s} a_r z_r > 0 \tag{33}
\]
By Lemma G.1 below, \( 1^T a_s = 1 \) and \( |a_s| \leq 1 \). Then \( (33) \) follows from Lemma G.2 below provided \( |s| \leq 8 \). Hence the
Jacobian is negative definite. The proof of Theorem 4.1 is complete after Lemmas [G.1] and [G.2] are proved.

To show that it satisfies C3, it follows directly from (27) that if \( x_i = 0 \) then \( \phi_i(x_a) = \infty \). It is also clear from (27) that the converse holds. This proves C3.

**Lemma G.1:** Fix any integer \( p \geq 1 \). Given any \( x \in \mathbb{R}^n \), define a vector \( a \) in \( \mathbb{R}^m \) as follows:

\[
a_i = \frac{2x_i}{\sum_{j=1}^{m} x_j} - \frac{x_i^p}{\sum_{j=1}^{m} x_j^p}, \quad 1 \leq i \leq m
\]

Then \( \sum_{i=1}^{m} a_i = 1 \) and \( \sum_{i=1}^{m} a_i^2 \leq 1 \).

**Proof:** It is obvious that \( \sum_{i=1}^{m} a_i = 1 \). To show \( \sum_{i=1}^{m} a_i^2 \leq 1 \), we have

\[
\sum_{i=1}^{m} a_i^2 = \frac{\sum_{i=1}^{m} x_i^{2p}}{(\sum_{j=1}^{m} x_j)^2} + \frac{4 \sum_{i=1}^{m} x_i^2}{(\sum_{j=1}^{m} x_j)^2} - \frac{4 \sum_{i=1}^{m} x_i^{p+1}}{(\sum_{j=1}^{m} x_j)^2} \left( \sum_{j=1}^{m} x_j \right)
\]

\[
\leq 1 + \frac{4 \sum_{i=1}^{m} x_i^2}{(\sum_{j=1}^{m} x_j)^2} - \frac{4 \sum_{i=1}^{m} x_i^{p+1}}{(\sum_{j=1}^{m} x_j)^2} \left( \sum_{j=1}^{m} x_j \right)
\]

\[
= 1 - 4 \frac{\sum_{i \leq j \leq m} x_i x_j (x_i - x_j) (x_i^{p-1} - x_j^{p-1})}{(\sum_{j=1}^{m} x_j)^2 (\sum_{j=1}^{m} x_j^p)}
\]

\[
\leq 1
\]

**Lemma G.2:** Let \( a \in \mathbb{R}^m \) that satisfies \( \sum_{i=1}^{m} a_i = 1 \) and \( \sum_{i=1}^{m} a_i^2 \leq 1 \). Then for any nonzero \( z \in \mathbb{R}^m \) we have

\[
f(z) := \sum_{i=1}^{m} z_i^2 + \sum_{i=1}^{m} z_i \sum_{i=1}^{m} a_i z_i > 0
\]

provided \( m \leq 8 \).

**Proof:** Given any \( M \) let \( Z_M := \{ z \mid \sum_{i=1}^{m} z_i = M \} \). It then suffices to show that, for every \( M \in \mathbb{R} \), \( f(z) > 0 \) for \( z \in Z_M \). Given any \( M \), consider

\[
\min_{z \in Z_M} f(z) = \min_{z \in Z_M} \sum_{i=1}^{m} z_i^2 + M \sum_{i=1}^{m} a_i z_i
\]

(35)

Its Lagrangian is

\[
L(z, \mu) = \sum_{i=1}^{m} z_i^2 + M \sum_{i=1}^{m} a_i z_i + \mu \left( \sum_{i=1}^{m} z_i - M \right)
\]

where \( \mu \) is the Lagrange multiplier. Setting \( \partial L / \partial z_i = 0 \) for all \( 1 \leq i \leq m \) and substitute it into \( \sum_{i=1}^{m} z_i = M \), we obtain the unique minimizer given by \( \mu = -3M/m \) and \( z_i = \frac{M}{2} (\frac{3}{m} - a_i) \). Then

\[
\min_{z \in Z_M} f(z) = \frac{M^2}{4} \left( \frac{9}{m} - \sum_{i=1}^{m} a_i^2 \right) \geq \frac{M^2}{4} \left( \frac{9}{m} - 1 \right)
\]

Hence, when \( M \neq 0 \), \( \min_{z \in Z_M} f(z) \geq 0 \) if \( n < 9 \). When \( z \) is nonzero but \( M = 0 \), then \( f(z) > 0 \) from (35). 

\[\text{3If } \beta = 0 \text{ the Jacobian degenerates to}
\]

\[
\frac{\partial \Phi_s}{\partial x_s} = -\frac{4}{\tau^2 ||x_s||^2} 11^T,
\]

which is merely negative semidefinite.

**APPENDIX H**

**PROOF OF LEMMA 3.1**

By the definition of \( D_k(A_k) \), we have

\[
\mathbb{E} \left[ D_k(A_k) \mid \sum_{i,j} a_{ij} \geq 1 \right] = d_k \mathbb{P} \left( \sum_j a_{kj} \geq 1 \mid \sum_{i,j} a_{ij} \geq 1 \right)
\]

\[
= d_k \mathbb{P} \left( \sum_j a_{kj} \geq 1 \right)
\]

\[
= d_k \left( \frac{q_k |A_k| \prod_i (1 - q_i) |A_i|}{\sum_i q_i |A_i|} \right) + o \left( \sum_i q_i \right)
\]

where the last equality follows from the independence of \( a_{ij} \) and \( \mathbb{P} \left( \sum_j a_{kj} \geq 1 \right) = 1 - (1 - q_k)^{|A_k|} = |A_k| q_k + o(q_k) \).

Thus,

\[
\mathbb{E} \left[ \sum_i D_k(A_k) \mid \sum_{i,j} a_{ij} \geq 1 \right] = \sum_k d_k q_k |A_k| + o \left( \sum_k q_k \right)
\]