ON THE CONVERGENCE OF A WEAK GREEDY ALGORITHM FOR THE MULTIVARIATE HAAR BASIS

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Abstract. We define a family of weak thresholding greedy algorithms for the multivariate Haar basis for $L_1[0,1]^d$ ($d \geq 1$). We prove convergence and uniform boundedness of the weak greedy approximants for all $f \in L_1[0,1]^d$.

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1. Introduction

Let \( \Psi = (\psi_n)_{n=1}^{\infty} \) be a semi-normalized Schauder basis for a Banach space \( X \). For \( f \in X \), let \( (c_n(f))_{n=1}^{\infty} \) denote the sequence of basis coefficients for \( f \). The Thresholding Greedy Algorithm (TGA) was introduced by Temlyakov [7] for the trigonometric system and subsequently extended to the Banach space setting by Konyagin and Temlyakov [5]. See [9] and the recent monograph [10] for the history of the problem and for background information on greedy approximation. The algorithm is defined as follows. For \( f \in X \) and \( n \geq 1 \), let \( \Lambda_n(f) \subset \mathbb{N} \) be the indices corresponding to a choice of \( n \) largest coefficients of \( f \) in absolute value, i.e., \( \Lambda_n(f) \) satisfies

\[
\min \{ |c_i(f)| : i \in \Lambda_n(f) \} \geq \max \{ |c_i(f)| : i / \in \Lambda_n(f) \}.
\]
(Note that \( \Lambda_n(f) \) is uniquely defined if and only if there is strict inequality in (1).) We call \( G_n(f) := \sum_{i \in \Lambda_n(f)} c_i(f) \psi_i \) an \( n \)th greedy approximant to \( f \) and say that the TGA converges if \( G_n(f) \to f \). The basis \( \Psi \) is said to be quasi-greedy if there exists \( K < \infty \) such that for all \( f \in X \) and \( n \geq 1 \), we have \( \|G_n(f)\| \leq K\|f\| \). Wojtaszczyk [11, Theorem 1] proved that \( \Psi \) is quasi-greedy if and only if the TGA converges for all initial vectors \( f \in X \).

It was proved in [3, Remark 6.3] that the one-dimensional Haar basis for \( L_1[0,1] \) (normalized in \( L_1[0,1] \)) is not quasi-greedy, i.e., that the TGA does not converge for certain initial vectors \( f \). However, it was proved in [4] that there is a weak thresholding greedy algorithm (WTGA) for the Haar basis which converges.

A WTGA is a procedure of the following general type. Fix a weakness parameter \( t \) with \( 0 < t < 1 \). For each \( f \in X \), define an increasing sequence \( (\Lambda_n^t(f)) \) of sets, consisting of of \( n \) coefficient indices, such that

\[
\min \{ |c_i(f)| : i \in \Lambda_n^t(f) \} \geq t \max \{ |c_i(f)| : i / \in \Lambda_n(f) \}.
\]

The WTGA is said to converge if the sequence of weak greedy approximants \( G_n^t(f) := \sum_{i \in \Lambda_n^t(f)} c_i(f) \psi_i \) converges to \( f \). It was proved in [6] that quasi-greediness of \( \Psi \) guarantees convergence for every WTGA.

However, if \( \Psi \) is not quasi-greedy then the index sets \( (\Lambda_n^t(f)) \) must be carefully chosen to ensure convergence. For the WTGA defined in
it was proved that the algorithm converges and that the weak greedy approximants are uniformly bounded, i.e., that \( \| G_n^t(f) \| \leq K(t) \| f \| \), where \( K(t) \) depends only on the weakness parameter.

In [8] it was proved for the multivariate Haar system, normalized in \( L_p[0,1]^d \), for \( d \geq 1 \) and \( 1 < p < \infty \), that

\[
\| f - G_n(f) \| \leq C(p,d) \sigma_n(f),
\]

where \( \sigma_n(f) \) denotes the error in the best \( n \)-term approximation to \( f \) in the \( L_p \) norm using the multivariate Haar system. As remarked above, convergence fails for \( p = 1 \). The goal of the present paper is to extend the results of [4] to the multivariate Haar system for \( L_1[0,1]^d \). The case \( d = 2 \) is especially interesting from the point of view of practical applications, and we refer the reader to [1] (and its references) for a nice exposition of the two-dimensional discrete Haar wavelet transform and its use in image compression.

Some serious obstacles have to be overcome in extending the one-dimensional results to higher dimensions. These difficulties are for the most part already present in the case \( d = 2 \). On the other hand, the passage from \( d = 2 \) to \( d \geq 3 \) is relatively straightforward.

The first obstacle in extending the one-dimensional algorithm of [4], which impeded progress on this problem for a considerable period, is that “the obvious generalization” fails to converge. Therefore, a more complicated algorithm is required, which depends on two parameters: the weakness parameter \( t \) and a second parameter \( s \), where \( 0 < t < s \).

An important feature of the algorithm, which it shares with the simpler one-dimensional algorithm, is that the weak greedy approximant is updated by applying a basic greedy step to the residual vector \( R_n^t(f) = f - G_n^t(f) \). The form of this greedy step ensures that the algorithm is branch-greedy in the sense of [2]. Roughly speaking, this means that the selection of the next coefficient in the basic greedy step depends only on the natural (finite) data set for weak thresholding consisting of all pairs

\[
\{(i, c_i(R_n(f))): \mid c_i(R_n(f)) \mid \geq t \max_{j \geq 1} \mid c_j(R_n(f)) \mid \}.
\]

Let us now describe the contents of the paper. In Section 2 we recall the definition of the multivariate Haar system, describe the weak
threshoding greedy algorithm alluded to in the title of the paper, and state the Main Theorem. The proof of the Main Theorem, which is presented in Section 7, uses two key lemmas which are proved in Sections 3–6.

The main result of Section 3 is the norm estimate Lemma 3.4. The results in Section 4 are based on the combinatorics of dyadic cubes. The main result of this section is the first key lemma, namely the norm estimate Lemma 4.6. Sections 5 and 6 are independent of Section 4. Section 5 contains an important symmetrization result. Section 6 is devoted to the second key lemma. The proof of this lemma uses an induction argument which makes essential use of the symmetrization results of the previous section.

In Section 8 we show that the algorithms diverge for the boundary cases $s = t$ and $s = 1$. This implies, in particular, that the multivariate Haar system is not quasi-greedy (Corollary 8.1).

## 2. Multivariate Haar system, Definition of the Algorithm, and Main Theorem

In this paper we consider greedy algorithms for the multivariate Haar system. Let $D_n$ be the set of dyadic intervals of length $2^{-n}$ and let $D^n_d := D_n \times \ldots \times D_n$ be the collection of all $d$-dimensional dyadic cubes of side length $2^{-n}$. Further, let $D := \bigcup_{n \geq 1} D_n$ and $D^d := \bigcup_{n \geq 1} D^n_d$.

For $a < b$, let

\[
\begin{align*}
    r^{(0)}_{[a,b)} &:= \frac{\chi_{[a,b)} - \chi_{b}}{b-a}, \\
    r^{(1)}_{[a,b)} &:= \frac{\chi_{[a,\frac{a+b}{2})} - \chi_{\frac{a+b}{2},b)}{b-a},
\end{align*}
\]

where $\chi$ denotes the characteristic function.

There are $2^d - 1$ different Haar functions corresponding to the dyadic cube $I \in D^d$, namely, for $1 \leq j \leq 2^d - 1$,

\[
h^{(j)}_I(x) := \prod_{k=1}^d r^{(\epsilon_k)}_{I_k}(x_k),
\]

where $x = (x_1, \ldots, x_d) \in [0,1)^d$, $I = I_1 \times \ldots \times I_d$ and $\epsilon_k \in \{0,1\}$ are defined from the binary representation $j = \sum_{k=1}^d \epsilon_k 2^{d-k}$. The Haar system is the set of all functions $h^{(j)}_I$ together with $\chi_{[0,1)^d}$. The Haar
coefficients are defined by
\[ c_I^{(i)}(f) = \mu(I) \int_I f h_I^{(i)}, \]
where the integral is taken with respect to \((d\text{-dimensional})\) Lebesgue measure \(\mu\). We write \([a, b) \prec [c, d)\) if \(b - a > d - c\) or \(b - a = d - c\) and \(a < c\). Further, for \(I = I_1 \times \ldots \times I_d \in \mathcal{D}^d, J = J_1 \times \ldots \times J_d \in \mathcal{D}^d\), we write \(I \prec J\) if \(I\) precedes \(J\) in the lexicographic ordering, i.e.,

- \(I_1 < J_1\), or
- \(I_1 = J_1\) and \(I_2 < J_2\), or

\[ \ldots \ldots \]
- \(I_1 = J_1, \ldots, I_{d-1} = J_{d-1}\) and \(I_d < J_d\).

Finally, we write \((I, i) < (J, j)\) if \(I \prec J\) or \(I = J\) and \(i < j\).

Note that each \(I \in \mathcal{D}_n^d\) is the disjoint union of \(2^d\) subcubes belonging to \(\mathcal{D}_{n+1}^d\) which we shall refer to as the immediate successors of \(I\).

Now we are ready to define the algorithm. For any \(f \in L_1[0, 1]^d\) and \(0 < t \leq s \leq 1\) we define the sequence \(\{G_{m}^{s,t}(f)\}\) inductively. We put
\[ G_{s,t}^0(f) = 0, R_{s,t}^0(f) = f \]
and for each \(m \geq 1\) we define \(G_{s,t}^m(f)\) and \(R_{s,t}^m(f)\) in the following way

1) Find the first cube in the order \(<\), denoted \(\Delta_m\), and then the smallest value of \(j_m\), with \(1 \leq j_m \leq 2^d - 1\), such that \((\Delta_m, j_m)\) satisfies
\[ |c_{\Delta_m}^{(j_m)}(R_{m-1}^{s,t})| = \max_{I \in \mathcal{D}^d} |c_I^{(i)}(R_{m-1}^{s,t})| \]

2) Define \(\tilde{\Delta}_m \supseteq \Delta_m\) to be the largest cube containing \(\Delta_m\) such that for every \(I \in \mathcal{D}^d\) with \(\Delta_m \subseteq I \subseteq \tilde{\Delta}_m\) there exists \(1 \leq i \leq 2^d - 1\) such that
\[ |c_I^{(i)}(R_{m-1}^{s,t})| \geq s \cdot |c_{\Delta_m}^{(j_m)}(R_{m-1}^{s,t})| \]

3) Find \(1 \leq i_m \leq 2^d - 1\) for which \(|c_{\Delta_m}^{(i_m)}(R_{m-1}^{s,t})|\) is the smallest value to satisfy
\[ |c_{\Delta_m}^{(i_m)}(R_{m-1}^{s,t})| \geq \frac{t}{s} \max_{1 \leq j \leq 2^d - 1} |c_{\Delta_m}^{(j)}(R_{m-1}^{s,t})|. \]

4) Let
\[ G_{m}^{s,t}(f) := G_{m-1}^{s,t} + c_{\Delta_m}^{(i_m)}(f) h_{\Delta_m}^{(i_m)}(f), R_{m}^{s,t} := f - G_{m}^{s,t}(f) \]
be the updated greedy approximant and residual.
Remark 2.1. Step 3) may be modified by replacing the selection condition (4) by the following condition:

\[
| c_{\Delta}^{(im)}(R_{m-1}^{s,t}) | \geq t | c_{\Delta}^{(jm)}(R_{m-1}^{s,t}) | .
\]

This defines a second weak greedy algorithm which has exactly the same convergence properties as the first algorithm.

Now we can state our main result.

**Main Theorem.** Let \(0 < t < s < 1\). Then the weak greedy algorithms defined above converge, i.e., \(f = \lim_{m \to \infty} G_{s,t}^{m}(f)\) for every \(f \in L_1[0, 1]^d\). Moreover, the greedy approximants are uniformly bounded, i.e., for all \(f \in L_1[0, 1]^d\) and for all \(m \geq 1\),

\[
\|G_{s,t}^{m}(f)\| \leq C(d, s, t)\|f\|.
\]

In Section 7 we show that \(G_{s,t}^{m}\) does not converge when \(s = t\) or \(s = 1\). From this, in particular from the case \(s = t = 1\), it follows that the multivariate Haar system is not a quasi-greedy basis in \(L_1[0, 1]^d\).

### 3. Norm Estimates by Expansion Coefficients

For any \(f \in L_1[0, 1]^d\), let

\[
sp(f) := \{ (\Delta, i) : c_{\Delta}^{(i)}(f) \neq 0 \},
\]

and let

\[
SP(f) := \{ \Delta : (\Delta, i) \in sp(f) \text{ for some } 1 \leq i \leq 2^d - 1 \}.
\]

Also, for any \(I \in \mathcal{D}^d\), let

\[
P_{\mathcal{I}} f := f - \sum_{\Delta \subseteq \mathcal{I}} \sum_{i=1}^{2^d-1} c_{\Delta}^{(i)}(f)h_{\Delta}^{(i)}.
\]

Note that \(P_{\mathcal{I}} f\) is constant on \(\mathcal{I}\) and it coincides with \(f\) outside of \(\mathcal{I}\). Recall, that the Haar system has the following monotonicity property:

\[
\|P_{\mathcal{I}} f\| \geq \|P_{\mathcal{J}} f\|, \text{ for any } \mathcal{I} \subseteq \mathcal{J}.
\]

Also, let us denote the norm of \(f \in L_1[0, 1]^d\) on the set \(\Delta\) by \(\|f\|_{\Delta}\), i.e.,

\[
\|f\|_{\Delta} := \int_{\Delta} |f|.
\]

Now we formulate two basic lemmas.
Lemma 3.1. Let $f \in L^1[0,1]^d$ and $I, J \in \mathcal{D}^d$ and $J \subseteq I$. Then
\[ \|f\|_I \geq |c^{(i)}_J(f)| \text{ for any } 1 \leq i \leq 2^d - 1. \]

Proof. By (3),
\[ |c^{(i)}_J(f)| = \mu(J) \left| \int_J fh^{(i)}_J \right| \leq \int_J |f| = \|f\|_J \leq \|f\|_I. \]
\[ \square \]

Lemma 3.2. Let $f \in L^1[0,1]^d$ and suppose $|c^{(i)}_\Delta(f)| \leq 1$ for all $\Delta \in \mathcal{D}^d$ and for all $1 \leq i \leq 2^d - 1$. Then, for any dyadic cube $I \in \mathcal{D}^d$, one has
\[ \|P_I f\|_I \leq 1. \]

Proof. Let $k$ be defined by $\mu(I) := 2^{-kd}$. Let the chain of dyadic cubes $\{\Delta_j\}_{j=0}^{k-1}$ be defined by the following conditions:
1) $I \subset \Delta_j$, 2) $\mu(\Delta_j) = 2^{-jd}$.

It is clear that $P_I f$ is constant on $I$ and that for any $x \in I$ one has
\[ P_I f(x) = c^{(0)}_{\Delta_0}(f) + \sum_{j=0}^{k-1} \sum_{i=1}^{2^d-1} c^{(i)}_{\Delta_j}(f) h^{(i)}_{\Delta_j}(x). \]

Taking into account the fact that $|h^{(i)}_{\Delta_j}(x)| = 2^{jd}$, we conclude that
\[ \|P_I f\|_I \leq 2^{-kd} \left( 1 + \sum_{j=0}^{k-1} (2^d - 1) \cdot 2^{jd} \right) = 1. \]
\[ \square \]

Lemma 3.3. Suppose $J \subset I$ and $\mu(J) = \frac{\mu(I)}{2^d}$, where $I, J \in \mathcal{D}^d$. Then
\[ \|f\|_{I \setminus J} \geq \begin{cases} \frac{1}{4} \left( \|P_I(f)\|_I - \|P_J(f)\|_J \right) \\ \frac{1}{4} \left( |c^{(i)}_I| - |c^{(j)}_I| \right) \\ \frac{1}{4} \left( \|P_I(f)\|_I - |c^{(i)}_I| \right). \end{cases} \]

For $d = 2$ the constant $1/4$ may be improved to $1/2$.

Proof. Here we prove the lemma (with the improved constant of $1/2$) only for the case $d = 2$. The case $d \geq 3$ is proved in Section 9.

Let us prove the first statement. The first inequality in (10) below follows from the monotonicity property of the Haar system, while the
second follows from the fact that \( P_I(f) \) and \( P_J(f) \) are constant on \( I \) and \( J \) respectively:

\[
\| f \|_{I \setminus J} \geq | \int_I f - \int_J f | \geq | \int_I P_I(f) - \int_J P_J(f) | \geq | \| P_I(f) \|_I - \| P_J(f) \|_J | .
\]

Let us prove the second and third statements of the lemma. Let \( \delta := \mu(I) \). Let us denote the value of \( P_I(f) \) on the cube \( I \) by \( H \) and the value of \( P_J(f) \) on the cube \( J \) by \( H_1 \).

Clearly,

\[
\| P_I(f) \|_I = \delta | H | ,
\]

and, since \( d = 2 \),

\[
\| P_J(f) \|_J = \frac{\delta | H_1 |}{4} .
\]

Note that for some choice of signs \( \epsilon_1, \epsilon_2, \epsilon_3 = \pm 1 \) one has

\[
H_1 = H + \frac{\epsilon_1 c^{(1)}_I + \epsilon_2 c^{(2)}_I + \epsilon_3 c^{(3)}_I}{\delta}.
\]

Let \( a_i := \epsilon_i c^{(i)}_I \). Then by the monotonicity property of the Haar system one has

\[
\| f \|_{I \setminus J} \geq \frac{1}{4} \left( | H \delta + a_1 - a_2 - a_3 | + | H \delta - a_1 + a_2 - a_3 | + | H \delta - a_1 - a_2 + a_3 | \right).
\]

Combining this inequality with (11) and using the triangle inequality one obtains the second and third statements of the lemma with 1/4 replaced by 1/2.

\[\square\]

**Lemma 3.4.** Let \( K \subset J \subset I \) with \( \mu(J) = \frac{\mu(I)}{2^d} \) and \( \mu(K) = \frac{\mu(J)}{2^d} \). Then

\[
\| f \|_{I \setminus K} \geq \frac{1}{8} \left| c^{(i)}_J - c^{(j)}_J \right| ,
\]

for all \( 1 \leq i, j \leq 2^d - 1 \).
Proof. Using Lemma 3.3 one gets

\[ \| f \|_{I \backslash J} \geq \frac{1}{2} \left( \left| \frac{c_I^{(i)} - \| P_I(f) \|_I}{4} \right| + \| P_I(f) \|_I - \| P_J(f) \|_J \right) \]

and

\[ \| f \|_{J \backslash K} \geq \left| \frac{c_J^{(j)} - \| P_J(f) \|_J}{4} \right| . \]

Hence

\[ \| f \|_{I \backslash K} = \| f \|_{I \backslash J} + \| f \|_{J \backslash K} \]

\[ \geq \left| \frac{c_I^{(i)} - \| P_I(f) \|_I}{8} \right| + \left| \frac{c_J^{(j)} - \| P_J(f) \|_J}{4} \right| \]

\[ \geq \left| \frac{c_I^{(i)} - c_J^{(j)}}{8} \right| . \]

Lemma 3.5. Let \( K \subseteq J \subseteq I \). Then

\[ \| f \|_{I \backslash K} \geq \left| \frac{c_I^{(i)} - c_J^{(j)}}{16} \right| . \]

Proof. If \( I = J \) then the lemma follows from Lemma 3.3. Suppose \( I \neq J \). Put \( \Delta_0 = I \) and inductively define the chain \( \{ \Delta_k \}_{k=0}^m \) of dyadic cubes as follows:

\[ \mathcal{K} \subset \Delta_{k+1} \subset \Delta_k, \quad \mu(\Delta_{k+1}) = \frac{\mu(\Delta_k)}{2d} \quad (0 \leq k < m). \]

Note that \( \Delta_m = \mathcal{K} \) and that \( J = \Delta_p \) for some \( 0 < p < s \). So by Lemma 3.4 we get

\[ \| f \|_{I \backslash K} \geq \| f \|_{\Delta_0 \backslash \Delta_{p+1}} \geq \frac{1}{2} \sum_{k=0}^{p-1} \| f \|_{\Delta_k \backslash \Delta_{k+2}} \]

\[ \geq \left| \frac{c_I^{(i)} - c_{\Delta_1}^{(j)}}{16} \right| + \sum_{k=1}^{p-2} \left| \frac{c_{\Delta_k}^{(1)} - c_{\Delta_{k+1}}^{(1)}}{16} \right| + \left| \frac{c_{\Delta_{p-1}}^{(1)} - c_J^{(j)}}{16} \right| \]

\[ \geq \left| \frac{c_I^{(i)} - c_J^{(j)}}{16} \right| . \]
4. Collections of Dyadic Cubes and the First Key Lemma

**Definition 4.1.** Let $\mathcal{J} \subset \mathcal{I}$. The *chain* $\mathcal{C}(\mathcal{I}, \mathcal{J})$ is the set of dyadic cubes $\mathcal{K}$ such that $\mathcal{J} \subseteq \mathcal{K} \subseteq \mathcal{I}$.

**Definition 4.2.** A finite set $\mathcal{R} \subset \mathcal{D}$ is called a *generalized chain* if there exists $\mathcal{I}_{\text{max}} \in \mathcal{R}$ such that for any $\mathcal{J} \in \mathcal{R}$ one has

a) $\mathcal{J} \subseteq \mathcal{I}_{\text{max}}$,

b) $\mathcal{C}(\mathcal{I}_{\text{max}}, \mathcal{J}) \subseteq \mathcal{R}$.

The cube $\mathcal{I}_{\text{max}} \in \mathcal{R}$ is called the *maximal cube* of $\mathcal{R}$. If $\mathcal{I}_{\text{max}} \neq [0,1)^d$ then the smallest cube which strictly contains $\mathcal{I}_{\text{max}}$ is called the *father* of $\mathcal{R}$, denoted $F(\mathcal{R})$.

The following lemma is the analogue of [4, Lemma 4]. We refer the reader to [4] for the proof.

**Lemma 4.3.** The union of two generalized chains $\mathcal{R}_1$ and $\mathcal{R}_2$ is a generalized chain if and only if either $\mathcal{R}_1 \cap \mathcal{R}_2 \neq \emptyset$ or $F(\mathcal{R}_1) \in \mathcal{R}_2$ or $F(\mathcal{R}_2) \in \mathcal{R}_1$. Then either $F(\mathcal{R}_1 \cup \mathcal{R}_2) = F(\mathcal{R}_1)$ or $F(\mathcal{R}_1 \cup \mathcal{R}_2) = F(\mathcal{R}_2)$.

Let us recall two more definitions from [4].

**Definition 4.4.** Let $\mathcal{S}$ be a finite subset of $\mathcal{D}$. We say that $\{\mathcal{R}_1, \ldots, \mathcal{R}_k\}$ is the *minimal generalized chain representation* (MGCR) of $\mathcal{P}$ if

a) $\mathcal{S} = \bigcup_{i=1}^{k} \mathcal{R}_i$,

b) $\mathcal{R}_i$ is a generalized chain for all $1 \leq i \leq k$,

c) $\mathcal{R}_i \cup \mathcal{R}_j$ is not a generalized chain for any $1 \leq i < j \leq k$.

It is shown in [4] that any set $\mathcal{S}$ has a unique MGCR.

**Definition 4.5.** Let $\mathcal{S}$ be a finite subset of $\mathcal{D}$ and $\mathcal{I} \in \mathcal{S}$. We say that $\mathcal{J} \in \mathcal{S}$ is a *son* of $\mathcal{I}$ with respect to $\mathcal{S}$ if $\mathcal{C}(\mathcal{I}, \mathcal{J}) \cap \mathcal{S} = \{\mathcal{I}, \mathcal{J}\}$.

The set of all sons of $\mathcal{I}$ with respect to $\mathcal{S}$ will be denoted by $\text{son}(\mathcal{I}, \mathcal{S})$.

For any $\mathcal{P} \subset \mathcal{S}$, put

\[ \text{son}(\mathcal{P}, \mathcal{S}) := \bigcup_{\mathcal{I} \in \mathcal{P}} \text{son}(\mathcal{I}, \mathcal{S}), \]

and

\[ \text{son}^{k+1}(\mathcal{P}, \mathcal{S}) := \text{son}(\text{son}^k(\mathcal{P}, \mathcal{S}), \mathcal{S}). \]
Finally, let us define the sets

- \( \Lambda_0(S) = \{ I \in S : \text{son}(I, S) = \emptyset \} \),
- \( \Lambda_1(S) = \{ I \in S : | \text{son}(I, S) | = 1 \} \),
- \( \Lambda_2(S) = \{ I \in S : | \text{son}(I, S) | \geq 2 \} \).

By induction on the cardinality of \( S \) (see [4, p. 56]), one has

\[
| \Lambda_2(S) | < | \Lambda_0(S) | .
\]

Now we are ready to formulate and prove the first key lemma.

**Lemma 4.6.** Let \( 0 < t < s < 1 \), and let \( p, q \in L_1[0, 1]^d \) be given, with \( \text{sp}(p) \) finite and \( [0, 1]^d \notin \text{SP}(p) \). Let \( \{ \mathcal{R}_1, \ldots, \mathcal{R}_k \} \) be the MGCR of \( \text{SP}(p) \). Suppose

1) \( \text{sp}(p) \cap \text{sp}(q) = \emptyset \),
2) for any \( \Delta \in \text{SP}(p) \) there exists \( 1 \leq i \leq 2^d - 1 \) such that \( | c^{(i)}_\Delta(p) | \geq s \),
3) if \( \Delta \in \text{SP}(p) \cap \text{SP}(q) \) then \( | c^{(i)}(q) \text{sp}(p) | < \frac{t}{s} | c^{(i)}(p) | \) for some \( (\Delta, i) \in \text{sp}(p) \) and for any \( (\Delta, j) \in \text{sp}(q) \),
4) for every \( 1 \leq l \leq k \) there exists \( \Delta \in \mathcal{R}_l \) and \( 1 \leq i \leq 2^d - 1 \) such that \( | c^{(i)}_{F(\mathcal{R}_l)}(q) | < s | c^{(i)}(p) | \) for all \( 1 \leq j \leq 2^d - 1 \).

Then

\[
\| p + q \| > C(s, t) | \mathcal{M} | ,
\]

where

\[
\mathcal{M} = \left( \bigcup_{i=1}^{k} \{ F(\mathcal{R}_l) \} \right) \cup \left( \text{SP}(p) \cap \text{SP}(q) \right).
\]

Note that \( [0, 1]^d \notin \text{SP}(p) \) by assumption, and so \( F(\mathcal{R}_1) \) exists for each \( 1 \leq l \leq k \).

**Proof.** We will proof the lemma for \( d = 2 \) (the extension to \( d \geq 3 \) is routine but the notation is more cumbersome). We consider two cases.

**CASE 1.** \( | \Lambda_0(\mathcal{M}) | \geq \frac{|\mathcal{M}|}{12} \). Let \( I \in \Lambda_0(\mathcal{M}) \). Then either a) \( I \in \text{SP}(p) \) or b) \( I = F(\mathcal{R}_l) \) for some \( 1 \leq l \leq k \).

a) If \( I \in \text{SP}(p) \) then combining Lemma 3.1 and condition 2) of this lemma, we conclude that

\[
\| p + q \|_I \geq s.
\]
b) If $\mathcal{I} = F(\mathcal{R}_I)$ then the maximal cube of $\mathcal{R}_I$ belongs to $SP(p)$ and therefore satisfies condition 2) of the Lemma. By the same argument as in the case a) we conclude that (17) holds for $\mathcal{I}$.

Note that the cubes $\mathcal{I} \in \Lambda_0(\mathcal{M})$ are disjoint, so

$$\|p + q\| \geq \frac{s}{12} |\mathcal{M}|.$$  

CASE 2. $|\Lambda_0(\mathcal{M})| < \frac{|\mathcal{M}|}{12}$. Taking into account (15) and the fact that $\Lambda_0(\mathcal{M})$ is always nonempty, one has

$$|\Lambda_1(\mathcal{M})| > \frac{|\mathcal{M}|}{6} + 1 \quad \text{and} \quad |\mathcal{M}| > 12.$$  

For $\mathcal{I} \in \Lambda_1(\mathcal{M})$, let $\mathcal{J}$ denote the unique son of $\mathcal{I}$. Let us define

$$(19) \quad \tilde{\Lambda}_1(\mathcal{M}) = \{ \mathcal{I} \in \Lambda_1(\mathcal{M}) : \mathcal{J} \in \Lambda_1(\mathcal{M}) \}.$$  

Note that

$$|\tilde{\Lambda}_1(\mathcal{M})| \geq |\Lambda_1(\mathcal{M})| - |\Lambda_0(\mathcal{M})| - |\Lambda_2(\mathcal{M})|$$

$$> \left( \frac{5|M|}{6} + 1 \right) - \left| \frac{|\mathcal{M}|}{12} - \frac{|\mathcal{M}|}{12} \right|$$

$$= \frac{2|\mathcal{M}|}{3} + 1.$$  

For $\mathcal{I} \in \tilde{\Lambda}_1(\mathcal{M})$, whose unique son (we recall) is denoted by $\mathcal{J}$, let $\mathcal{K}$ denote the unique son of $\mathcal{J}$. Let us prove that either

$$(21) \quad \|p + q\|_{\mathcal{I}\setminus \mathcal{J}} \geq C_1(s, t)$$

or

$$(22) \quad \|p + q\|_{\mathcal{J}\setminus \mathcal{K}} \geq C_1(s, t).$$

If $\mathcal{I} \in SP(p) \cap SP(q)$ then by conditions 2) and 3) of the lemma, we have $|c^{(i)}_\mathcal{I}(p)| \geq s$ and $|c^{(j)}_\mathcal{I}(q)| < \frac{t}{s}$ for some $i$ and $j$. Therefore, according to Lemma 3.5, we have

$$\|p + q\|_{\mathcal{I}\setminus \mathcal{J}} \geq \frac{s-t}{16}.$$  

Using the same argument for the case when $\mathcal{J} \in SP(p) \cap SP(q)$, we have

$$\|p + q\|_{\mathcal{J}\setminus \mathcal{K}} \geq \frac{s-t}{16}.$$
It remains to consider the case when $\mathcal{I}, \mathcal{J} \notin SP(p) \cap SP(q)$. Then we have

$$\mathcal{I} = F(\mathcal{R}_{l_1}) \text{ and } \mathcal{J} = F(\mathcal{R}_{l_2}),$$

for some $1 \leq l_1, l_2 \leq k$. Let $\Delta$ and $i$ be chosen according to conditions 2) and 4) of the Lemma for the generalized chain $\mathcal{R}_{l_1}$. Then we have that

\begin{equation}
| c^{(i)}_\Delta(p) | \geq s, \tag{23}
\end{equation}

and

\begin{equation}
| c^{(1)}_\mathcal{I}(q) | < s | c^{(i)}_\Delta(p) |. \tag{24}
\end{equation}

By Lemma 4.3 $\mathcal{R}_{l_1}$ and $\mathcal{R}_{l_2}$ are disjoint, so $\mathcal{J} \notin \mathcal{R}_{l_1}$ and hence $\Delta \not\subseteq \mathcal{J}$ by the definition of generalized chain. In the case when $\Delta \cap \mathcal{J} = \emptyset$ we get

$$\| p + q \|_{\mathcal{I} \setminus \mathcal{J}} \geq s$$

directly from (23) and Lemma 3.1. Finally, for the case $\mathcal{J} \subsetneq \Delta$ we apply (23), (24) and Lemma 3.5 to conclude that

$$\| p + q \|_{\mathcal{I} \setminus \mathcal{J}} \geq s(1 - s) \frac{1}{16}.$$ 

So, for every $\mathcal{I} \in \Lambda_1(M)$, either (21) or (22) holds. Recall that $\mathcal{J}$ and $\mathcal{K}$ depend on $\mathcal{I}$. It is to easily seen that the sets $\mathcal{I} \setminus \mathcal{J}$ as $\mathcal{I}$ ranges over $\Lambda_1(M)$ are disjoint. The same is true for the sets $\mathcal{J} \setminus \mathcal{K}$ as $\mathcal{I}$ ranges over $\tilde{\Lambda}_1(M)$. Therefore

$$\| p + q \| \geq \sum_{\mathcal{I} \in \Lambda_1(M)} \| p + q \|_{\mathcal{I} \setminus \mathcal{J}},$$

and

$$\| p + q \| \geq \sum_{\mathcal{I} \in \tilde{\Lambda}_1(M)} \| p + q \|_{\mathcal{J} \setminus \mathcal{K}}.$$ 

Using (20) we conclude that

$$\| p + q \| \geq \frac{1}{2} \sum_{\mathcal{I} \in \Lambda_1(M)} \left( \| p + q \|_{\mathcal{I} \setminus \mathcal{J}} + \| p + q \|_{\mathcal{J} \setminus \mathcal{K}} \right) \geq C(s, t) | M |.$$ 

$\square$
5. Symmetrization Properties

We will prove the results in this section only for $d = 2$, but the extension to $d \geq 3$ is routine.

Let $\Delta = [a, a + 2\delta] \times [b, b + 2\delta)$ be a dyadic square of side length $2\delta$, and let $\Delta_i, 1 \leq i \leq 4$, be the four disjoint immediate successor squares of $\Delta$ of side length $\delta$. For each $1 \leq i \leq 4$, let us denote by $L_i(f, \Delta)$ the function which agrees with $f$ on the sets $[0, 1)^2 \setminus \Delta$ and $\Delta_i$, and which

‘copies’ $f$ from the square $\Delta_i$ to the other three squares $\Delta_j, j \neq i$. More precisely, let $\Delta_j = u_j + \Delta_i$, where $u_j \in \mathbb{R}^2$. Then

$L_i(f, \Delta)(x) = f(x - u_j)$ for $1 \leq j \leq 4$ and for all $x \in \Delta_j$.

**Lemma 5.1.** Let $f, g \in L_1[0, 1]^2$, let $SP(f + g) \neq \emptyset$, and let

(25) $B := \frac{\|f\|}{\|f + g\|}$.

Then (for any square $\Delta$) we have either

(26) $\|L_i(f, \Delta)\| > B\|L_i(f + g, \Delta)\|$, for some $1 \leq i \leq 4$,

or

(27) $\|L_i(f, \Delta)\| = B\|L_i(f + g, \Delta)\|$, for all $1 \leq i \leq 4$.

**Proof.** Let us assume that the statement of the lemma is not correct. Then we have

(28) $\sum_{i=1}^{4} \|L_i(f, \Delta)\| < B \sum_{i=1}^{4} \|L_i(f + g, \Delta)\|$.

Note that

$\|L_i(f, \Delta)\| = \|f\| + 4\|f\|\Delta_i - \sum_{j=1}^{4} \|f\|\Delta_j$

and

$\|L_i(f + g, \Delta)\| = \|f + g\| + 4\|f + g\|\Delta_i - \sum_{j=1}^{4} \|f + g\|\Delta_j$.

By substituting the last two inequalities into (28), we conclude that,

$4\|f\| < 4B\|f + g\|$.
which is contradiction.

\[ \square \]

**Remark 5.2.** If \( f \neq 0 \) then \( L_j(f, \Delta) \neq 0 \) for some \( 1 \leq j \leq 4 \).

For \( f, g \in L_1[0, 1]^2 \), with \( f + g \neq 0 \), and \( \Delta \in D^2 \), define
\begin{align*}
L((f, g), \Delta) := \begin{cases} 
(L_i(f, \Delta), L_i(g, \Delta)), & \text{if } (26) \text{ holds and } \|L_i(f + g, \Delta)\| > 0 \\
(L_j(f, \Delta), L_j(g, \Delta)), & \text{where } \|L_j(f + g, \Delta)\| > 0 \text{ otherwise.}
\end{cases}
\end{align*}

**Lemma 5.3.** Let \( f, g \in L_1[0, 1]^2 \) and \( \Delta \in D^2 \) satisfy
\begin{enumerate}
\item \( \Delta \notin SP(f), \Delta \notin SP(g) \),
\item \( sp(f) \cap sp(g) = \emptyset \),
\item \( sp(f + g) \neq \emptyset \).
\end{enumerate}

Further, let \( (f', g') := L((f, g), \Delta) = L_i((f, g), \Delta) \) for some \( 1 \leq i \leq 4 \). Then,
\begin{enumerate}
\item \( \Delta \notin SP(f'), \Delta \notin SP(g') \),
\item \( c_T^{(j)}(f') = c_T^{(j)}(f) \) and \( c_T^{(j)}(g') = c_T^{(j)}(g) \) for all \( I \) with \( I \not\subset \Delta \setminus \Delta_i \),
\item for every \( I \subset \Delta \setminus \Delta_i \) there exists \( J \in \Delta_i \) such that \( c_T^{(j)}(f') = c_T^{(j)}(f) \) and \( c_T^{(j)}(g') = c_T^{(j)}(g) \) for all \( 1 \leq j \leq 4 \),
\item \( sp(f') \cap sp(g') = \emptyset \),
\item \( sp(f' + g') \neq \emptyset \),
\item \( \frac{\|f'\|}{\|f' + g'\|} \geq \frac{\|f\|}{\|f + g\|} \).
\end{enumerate}

**Proof.** From condition i) of the lemma and (3) it follows that
\begin{align*}
(30) \quad \int_{\Delta_j} f = \int_{\Delta_k} f \quad \text{and} \quad \int_{\Delta_j} g = \int_{\Delta_k} g \quad \text{for all } 1 \leq j, k \leq 4.
\end{align*}

The functions \( f' \) and \( g' \) (replacing \( f \) and \( g \)) will also satisfy (30), which gives statement 1). Clearly, \( \int_{\Delta_i} f = \int_{\Delta_i} f' \) and \( \int_{\Delta_i} g = \int_{\Delta_i} g' \), and hence
\begin{align*}
(31) \quad \int_{\Delta} f = \int_{\Delta} f' \quad \text{and} \quad \int_{\Delta} g = \int_{\Delta} g'.
\end{align*}

If \( I \not\subset \Delta \setminus \Delta_i \), then either \( I \cap \left( \Delta \setminus \Delta_i \right) = \emptyset \) or \( I = \Delta \) or \( I \supseteq \Delta \). In the first case we have \( f = f' \) and \( g = g' \) on \( I \) and therefore we have statement 2). The second case is equivalent to statement 1). In the last case it is easy to check that \( c_T^{(j)}(f') \) and \( c_T^{(j)}(g') \) depend only on the average values \( \int_{\Delta} f' \) and \( \int_{\Delta} g' \) on \( \Delta \), which implies, taking (31) into
account, that \( c^{(j)}_I(f') = c^{(j)}_I(f) \) and \( c^{(j)}_I(g') = c^{(j)}_I(g) \). Thus, the proof of statement 2) is complete. Statement 3) is obvious and follows from the definition of the operator \( L \).

Statement 4) follows from the condition ii) and statements 2) and 3). Statement 5) follows immediately from the definition of \( (L(f, g), \Delta) \).

Finally, let us prove statement 6). The statement is obvious if \( f' \) and \( g' \) are defined according to the first line of (29). Let us consider the case when they are defined according to the second line. We have two cases.

CASE 1. Suppose (26) fails. Then, by Lemma 5.1, the set of equalities (27) holds, which implies that we have equality in statement 6).

CASE 2. Now suppose (26) holds, but that \( L_i(f + g, \Delta) = 0 \). Since \( sp(L_i(f)) \cap sp(L_i(g)) = \emptyset \) we conclude that \( L_i(f) = L_i(g) = 0 \). So inequality in (26) is impossible, which contradicts our assumption. □

6. The Second Key Lemma

**Lemma 6.1.** Let \( f, g \in L_1[0, 1]^d \), with \( sp(f) \) finite, and let \( \{R_1, \ldots, R_k\} \) be the MGCR of \( SP(f) \). Suppose that

1) \( SP(f) \cap SP(g) = \emptyset \),
2) \([0,1]^d \not\in SP(f)\) and \([0,1]^d \not\in SP(g)\),
3) \( F(R_l) \not\in SP(g) \) for all \( 1 \leq l \leq k \),
4) \( |c^{(j)}_I(g)| \leq 1 \) for all \( I \in \mathcal{D} \) and \( 1 \leq j \leq 2^d - 1 \),
5) for every \( 1 \leq l \leq k \) there exist \( I \in R_l \) and \( 1 \leq j \leq 2^d - 1 \) such that \( |c^{(j)}_I(f)| \geq t \).

Then

\[
\frac{\|f\|}{\|f + g\|} \leq C(t).
\]

**Proof.** We may assume that \( \mu(F(R_1)) \leq \mu(F(R_2)) \leq \ldots \leq \mu(F(R_k)) \). Let

\[
(f_1, g_1) := L(f, g, F(R_1)).
\]

Let \( V_1 \subset F(R_1) \) denote the cube from which the values of \( f \) and \( g \) are copied to the other immediate successor cubes of \( F(R_1) \) to define \( L(f, g) \). Note that \( SP(f_1) \) is obtained as follows:

1) ‘Remove’ all generalized chains from \( MGCR(f) \) whose maximal cubes are contained in \( F(R_1) \setminus V_1 \),
2) ‘Copy’ all generalized chains from $MGCR(f)$ whose maximal cubes are contained in $V_1$ to the other $2^d - 1$ immediate successor cubes of $F(R_1)$.

Then $f_1$ and $g_1$ have all properties that are listed in Lemma 5.3. In particular,

$$\frac{\|f\|}{\|f + g\|} \leq \frac{\|f_1\|}{\|f_1 + g_1\|}.$$ 

Inductively, for $i = 2, \ldots, k$, let $$(f_i, g_i) := L(f_{i-1}, g_{i-1}, F(R_i)).$$

Finally, let $f' := f_k$ and $g' := g_k$. It is easy to check that $f'$ and $g'$ have the following properties:

P1) $SP(f') \cap SP(g') = \emptyset$,

P2) $[0, 1)^d \notin SP(f')$ and $[0, 1)^d \notin SP(g')$,

P3) $F(R_i') \notin SP(g')$ for all $1 \leq l \leq k'$,

P4) $|c^{(j)}_{(I)}(f')| \leq 1$ for all $I \in D$ and $1 \leq j \leq 2^d - 1$,

P5) for every $1 \leq l \leq k'$ there exists $I \in R_i'$ and $1 \leq j \leq 2^d - 1$ such that $|c^{(j)}_{(I)}(f')| \geq t$,

P6) for every $1 \leq l \leq k'$ the functions $f'$ and $g'$ are ‘copied’ from one immediate successor cube of the cube $F(R_i')$ to all the other $2^d - 1$ immediate successors,

P7) $$\frac{\|f\|}{\|f + g\|} \leq \frac{\|f'\|}{\|f' + g'\|},$$

where $\{R'_1, \ldots, R'_{k'}\}$ is the MGCR of $SP(f')$. Let

$$S := \{F(R'_i) : 1 \leq l \leq k'\}.$$ 

Suppose that $c^{(j)}_{(I)}(f') \neq 0$. Then $I$ belongs to some generalized chain $R'_i$ and therefore there exists $J \in S$ such that $I \subset J$. So we conclude that

$$supp(f') \subseteq \bigcup_{I \in S} I.$$ 

We say that $I \in S$ has order $k$ if

$$son_k(I, S) \neq \emptyset \quad \text{and} \quad son^{k+1}(I, S) = \emptyset.$$ 

Let us prove that for any $I \in S$

$$\|f'\|_I < (5t^{-1} + 2)\|f' + g'\|_I - 2t - 8.$$
We use induction on the order of $I$. Suppose the order of $I$ is 0. Let $I_1, I_2, \ldots, I_{2^d}$ be the immediate successor cubes of $I$. Since $I \in S$ then, taking into account P5), P6) and Lemma 3.1, one has

$$\|f' + g'\|_{I_j} \geq t \text{ for all } 1 \leq j \leq 2^d.$$  

Hence

$$\|f' + g'\|_I \geq 2^d t. \tag{35}$$

It follows from Lemma 3.2 and P4) that

$$\|P_I(g')\|_I \leq 1. \tag{36}$$

Since the order of $I$ is equal to 0, the monotonicity of the Haar system and P1) give

$$\|f' + g'\|_I \geq \|f' + P_I(g')\|_I. \tag{37}$$

Combining (35), (36) and (37) gives

$$(5t^{-1} + 2) \|f' + g'\|_I \geq (5t^{-1} + 1) \|f' + g'\|_I + \|f' + P_I(g')\|_I \geq 2^d t(5t^{-1} + 1) + \|f' + P_I(g')\|_I \geq 10 + 2t + \|f' + P_I(g')\|_I \geq 2^d t + 2t + 9, \tag{38}$$

which gives (34) for $I$.

Assume now that (34) holds for all cubes of order $\leq k$. We will prove the estimate for all cubes $I \in S$ of order $k + 1$. Let $I_1, \ldots, I_{2^d}$ be the immediate successor cubes of $I$ and let $\text{son}(I, S) = \{J_1, J_2, \ldots, J_i\}$. Because of the symmetry property P6), each cube $I$ contains the same number of the $J_p$ cubes, $a$, say, where $a \geq 1$. Let $J_1, \ldots, J_a$ be contained in $I_1$. The cubes $J_i$ are disjoint and their orders are $\leq k$. Therefore, by the induction hypothesis,

$$\|f'\|_{J_i} \leq (5t^{-1} + 2) \|f' + g'\|_{J_i} - 2t - 8 \text{ for all } i = 1, 2, \ldots, a. \tag{39}$$

Let $D := I_1 \setminus \left( \bigcup_{i=1}^a J_i \right)$,

$$\alpha := P_{J_1}(P_{J_2}(\ldots P_{J_a}(f')\ldots)),$$

and

$$\beta := P_{J_1}(P_{J_2}(\ldots P_{J_a}(g')\ldots)).$$
For \( \alpha \) and \( \beta \), we have
\[
\alpha = f' \quad \text{and} \quad \beta = g' \quad \text{on } D,
\]
and by Lemma 3.2 and P4)
\[
\|\beta\|_{J_i} \leq 1 \quad \text{for any } 1 \leq i \leq s.
\]
By Lemma 3.2 and the monotonicity of the Haar system, we have
\[
\|\alpha + \beta\|_{I_1} \geq \|\alpha + \mathcal{P}_{I_1}(\beta)\|_{I_1} \geq \|\alpha\|_{I_1} - 1.
\]
Using (41), we get
\[
\|\alpha + \beta\|_{I_1 \setminus D} \leq \|\alpha\|_{I_1 \setminus D} + \|\beta\|_{I_1 \setminus D}
\]
\[
= \|\alpha\|_{I_1 \setminus D} + \sum_{i=1}^{a} \|\beta\|_{J_i}
\]
\[
\leq \|\alpha\|_{I_1 \setminus D} + a.
\]
From (40) and (42), we get
\[
\|f' + g'\|_{D} = \|\alpha + \beta\|_{D}
\]
\[
= \|\alpha + \beta\|_{I_1} - \|\alpha + \beta\|_{I_1 \setminus D}
\]
\[
\geq \|\alpha\|_{I_1} - 1 - \|\alpha\|_{I_1 \setminus D} - a
\]
\[
= \|\alpha\|_{D} - a - 1
\]
\[
= \|f'\|_{D} - a - 1.
\]
Combining this with (39) gives
\[
\|f'\|_{I_1} = \|f'\|_{D} + \sum_{i=1}^{a} \|f'\|_{J_i}
\]
\[
\leq \|f' + g'\|_{D} + a + 1 + (5t^{-1} + 2) \sum_{i=1}^{a} \|f' + g'\|_{J_i} - a(2t + 8)
\]
\[
\leq (5t^{-1} + 2)\|f' + g'\|_{I_1} - 2t - 6.
\]
Using the symmetry property P6), we conclude
\[
\|f'\|_{I} = 2^d\|f'\|_{I_1}
\]
\[
\leq (5t^{-1} + 2)\|f' + g'\|_{I} - 2^d(2t + 6)
\]
\[
< (5t^{-1} + 2)\|f' + g'\|_{I} - 2t - 8.
\]
So (31) holds for \( \mathcal{I} \). This completes the induction proof.
Let $B := \bigcup_{I \in S} I$. We can represent the set $B$ as the union of some disjoint intervals from $S$. Hence

$$\|f\|_B < (5t^{-1} + 2)\|f + g\|_B.$$ 

By (33) we have $f' = 0$ on the set $[0,1]^d \setminus B$, so

$$\|f\| < (5t^{-1} + 2)\|f + g\|.$$ 

This, together with P7), completes the proof. \hfill \Box

7. Main Results

Proof of the Main Theorem. Convergence is obvious if $f - G_m^{(s,t)}(f) = 0$. So assume that $G_m^{(s,t)}(f) \neq f$. Let us define

$$p := \frac{G_m^{(s,t)}(f)}{\max |c_{\Delta}^{(j)}(f - G_m^{(s,t)}(f)) : (\Delta, j) \in sp(f)|},$$

$$q := \frac{f - G_m^{(s,t)}(f)}{\max |c_{\Delta}^{(j)}(f - G_m^{(s,t)}(f)) : (\Delta, j) \in sp(f)|}.$$ 

It is clear that $p$ and $q$ satisfy all the conditions of Lemma 4.6. Therefore we have the estimate (16). Define

$$\tilde{q} := q - \sum_{I \in M} \sum_{j=1}^{2^d-1} c_{I}^{(j)}(q) h_{I}^{(j)}.$$ 

Clearly,

$$\|q - \tilde{q}\| \leq (2^d - 1) |M|.$$ 

Combining this with inequality (16) yields

$$\frac{\|p + \tilde{q}\|}{\|p + q\|} \leq 1 + \frac{\|q - \tilde{q}\|}{\|p + q\|} \leq 1 + \frac{2^d - 1}{C(s,t)}.$$ 

It remains to observe that functions $p$ and $\tilde{q}$ (in place of $f$ and $g$, respectively) satisfy the conditions of Lemma 5.3. Hence

$$\frac{\|p\|}{\|p + \tilde{q}\|} \leq C(t).$$ 

Combining this with (45) and (43) we get

$$\frac{\|G_m^{(s,t)}(f)\|}{\|f\|} \leq C(t) \cdot \left(1 + \frac{2^d - 1}{C(s,t)}\right),$$
which gives the uniform boundedness inequality \((6)\) of the Main Theorem with
\[
C(s, t, d) = C(t) \cdot \left(1 + \frac{2d - 1}{C(s, t)}\right).
\]
To deduce convergence of the algorithm from the uniform boundedness of the greedy approximants, we follow the argument from [11] for the Thresholding Greedy Algorithm. Let \(\Psi := (\psi_n)_{n=0}^\infty\) be the enumeration of the multivariate Haar basis in the order determined by \(\prec\). Then \(\Psi\) is a Schauder basis for \(L_1[0, 1]^d\). For \(N \geq 0\), let \(P_N\) denote the basis projection onto \(\text{span}(\psi_j)_{j=0}^N\). Suppose that \(N = (2d - 1)N_1 + 1\) so that \(P_N\) projects onto the span of all Haar functions supported on the first \(N_1\) dyadic cubes in the order \(\prec\). It is easily seen from the definition of the algorithm that for all sufficiently large \(n\) there exists \(m\) such that
\[
G_{s,t}^n(f) = P_N(f) + G_{s,t}^m(f - P_N(f)).
\]
Since \(\|G_{s,t}^m(f - P_N(f))\| \leq C(s, t, d)\|f - P_N(f)\| \to 0\) as \(N \to \infty\), we get \(f = \lim_{n \to \infty} G_{s,t}^n(f)\) as required. \(\square\)

**Corollary 7.1.** For every \(0 < t < 1\) there is a convergent implementation of the weak greedy algorithm with weakness parameter \(t\) for the multivariate Haar basis for \(L_1[0, 1]^d\) such that, for all \(f \in L_1[0, 1]^d\), we have
\[
\|G_{n}^t(f)\| \leq \frac{C_d}{1 - t}\|f\|,
\]
where \(C_d \ll 2^d\).

**Proof.** From the proof of Lemma 4.6, \(C(s, t) \geq \min(s(1 - s), s - t)/24\), and from the proof of Lemma 5.3, \(C(t) \leq 5/t + 12\). Setting \(s = (1 + t)/2\) and substituting these estimates into (16) gives the result for \(G_{n}^t(f) := G_{s,t}^n(f)\). \(\square\)

Finally, let us show that \(C_d \gg \sqrt{d}\). Let \((r_n)_{n=1}^\infty\) be the usual Rademacher functions defined on \([0, 1]\). For each finite \(A \subset \mathbb{N}\), recall that the Walsh function \(w_A := \prod_{n \in A} r_n\). The Walsh system \(W := \{w_A: A \in \mathbb{N}^{(\leq \infty)}\}\) is a fundamental orthogonal system for \(L_1[0, 1]\).

**Theorem 7.2.** Let \(0 < t < 1\). The greedy approximants with respect to \(W\) are unbounded for every implementation of the weak thresholding greedy algorithm with weakness parameter \(t\).
Proof. Let $0 < u < t$. For $N \geq 1$, consider
\[ f_N := \prod_{n=1}^{N} (1 + ur_n). \]
From the independence of the Rademacher functions and the fact that
$1 + ur_n \geq 0$, we get $\|f_N\| = 1$. Also, for any implementation of the weak
thresholding greedy algorithm with weakness parameter $t$, we have
\[ G_{N+1}^t(f_N) = 1 + u \sum_{n=1}^{N} r_n. \]
So, by Khinchine’s inequality,
\[ \|G_{N+1}^t(f_N)\| \geq u\| \sum_{n=1}^{N} r_n\| - 1 \gg u\sqrt{N}, \]
which gives the unboundedness of the greedy approximants. \qed

Corollary 7.3. $C_d \gg \sqrt{d}$.

Proof. It suffices to observe that the Haar functions $\chi_{[0,1)^d} \cup (h_{[0,1)^d})^{2d-1}$ have the same joint distribution as the initial segment of the Walsh
system $\{w_A: A \subset \{1, \ldots, d\}\}$. Hence the result follows from (47). \qed

8. The Boundary Cases $s = 1$ and $s = t$

In this section we will assume that $d = 2$, that $N = 2k$ is even
number, and that $0 < \epsilon < 1$. Let $\Delta_n := [0, \frac{1}{2^n}) \times [0, \frac{1}{2^n})$. Our starting
point is the function
\[ f_N = 1 + \sum_{n=0}^{N-1} \sum_{j=1}^{3} h_{\Delta_n}^{(j)}. \]
It is easy to check, that
\[ f_N = \begin{cases} 2^{2N} \text{ on } \Delta_N, \\ 0 \text{ otherwise.} \end{cases} \]
So $\|f_N\| = 1$.

Now let us consider the case when $s = 1$. Let
\[ f_N^\epsilon := 1 + \sum_{n=0}^{k-1} \sum_{j=1}^{3} (h_{\Delta_{2n+1}}^{(j)} + (1 - \epsilon)h_{\Delta_{2n}}^{(j)}). \]
It is clear that \( \| f_N \| \leq \| x_N \| + 3k\epsilon \). On the other hand, it is easy to check that

\[
G^{1,t}_{3k+1}(f_N^e) = 1 + \sum_{n=0}^{k-1} \sum_{j=1}^{3} h^{(j)}_{\Delta_{2n+1}}.
\]

By using Lemma 3.4 for \( I = \Delta_{2n+1}, J = \Delta_{2n+2}, \) and \( K = \Delta_{2n+3}, \) we get

\[
\| G^{1,t}_{3k+1}(f_N^e) \| \geq \sum_{n=0}^{k-1} \| G^{1,t}_{3k+1}(f_N^e) \|_{\Delta_{2n+1}\backslash\Delta_{2n+3}} \geq \frac{k}{8},
\]

so

\[
\frac{\| G^{1,t}_{3k+1}(f_N^e) \|}{\| f_N^e \|} \geq \frac{k}{8(1 + 3k\epsilon)}.
\]

Since \( k \) and \( \epsilon \) are arbitrary, we conclude that the operator \( G^{1,t}_{n} \) is not bounded.

Now, let us consider the case \( s = t \). Let

\[
g_N^e := t \left( 1 + \sum_{n=0}^{N-1} \left( h^{(1)}_{\Delta_n} + h^{(2)}_{\Delta_n} + (1 - \epsilon)h^{(3)}_{\Delta_n} \right) \right) + h^{(1)}_{\Delta_N}.
\]

Note that

\[
g_N^e = tf_N - tf_N^e + \sum_{n=0}^{N-1} h^{(3)}_{\Delta_n} + h^{(1)}_{\Delta_N},
\]

whence by the triangle inequality \( \| g_N^e \| \leq 1 + t + Nt\epsilon \). Note that

\[
G^{t,t}_{2N+1}(g_N^e) = t \left( 1 + \sum_{n=0}^{N-1} \left( h^{(1)}_{\Delta_n} + h^{(2)}_{\Delta_n} \right) \right).
\]

Hence by Lemma 3.3

\[
\| G^{t,t}_{2N+1}(g_N^e) \| \geq \sum_{n=0}^{N-1} \| G^{t,t}_{2N+1}(g_N^e) \|_{\Delta_n \backslash \Delta_{n+1}} \geq \frac{Nt}{2}.
\]

So

\[
\frac{\| G^{t,t}_{2N+1}(g_N^e) \|}{\| g_N^e \|} \geq \frac{Nt}{2(1 + t + Nt\epsilon)}.
\]

Since \( N \) and \( \epsilon \) are arbitrary, we conclude that the operator \( G^{t,t}_{n} \) is not bounded.

By a “gliding hump” argument (see [11] for the details) there exists \( f \in L_1[0, 1]^d \) for which \( (G^{t,t}_{n}(f)) \) diverges.

The case \( s = t = 1 \) implies the following result.
Corollary 8.1. The multivariate Haar system is not a quasi-greedy basis of $L_1[0,1]^d$.

9. Appendix: Proof of Lemma 3.3

Here we prove Lemma 3.3 for $d \geq 3$. The first statement is identical to the case $d = 2$. We prove the second and third statements. Let $H$ be the value of $S_I(f)$ on $I$ and also $\mu(I) = \delta$. Define signs $\sigma_p = \pm 1$ in such a way that the value of $S_{J}(f)$ is equal to $H + \frac{1}{\delta} \sum_{p=1}^{2^d-1} a_p$ where $a_p = \sigma_p c_I^{(p)}(f)$. Let $(J_j)_{j=1}^{2^d-1}$ be an enumeration of the $2^d - 1$ immediate successors of $I$ excluding $J$. For $1 \leq j \leq 2^d - 1$, let $\epsilon_j^{(p)} = \pm 1$ denote the value taken by $\delta \sigma_h^{(p)}$ on the cube $J_j$.

Then by monotonicity

\[ \| f \|_{I \setminus J} \geq \frac{\delta}{2^d} \sum_{j=1}^{2^d-1} | H + \frac{1}{\delta} \sum_{p=1}^{2^d-1} \epsilon_j^{(p)} a_p |. \]  

For a fixed $j$ (resp., for a fixed $p$) there are exactly $2^d - 1$ of the coefficients $\{ \epsilon_j^{(p)} \}$ that are equal to $-1$, and for distinct $p$ and $q$ the following orthogonality property is easily verified:

\[ \sum_{\{j: \epsilon_j^{(p)} = -1\}} \epsilon_j^{(q)} = 0. \]  

Let us fix some $p_0$. Using (53) we get

\[ \| f \|_{I \setminus J} \geq \frac{\delta}{2^d} \sum_{\{j: \epsilon_j^{(p_0)} = -1\}} \left( H + \frac{1}{\delta} \sum_{p=1}^{2^d-1} \epsilon_j^{(p)} a_p \right) \]

\[ = \frac{1}{2^d} \left| 2^{d-1} H \delta - 2^{d-1} a_{p_0} \right| = \left| \frac{H \delta - a_{p_0}}{2} \right|, \]

which proves the second statement of the lemma.

Finally, for distinct $p$ and $q$, we have

\[ \| f \|_{I \setminus J} \geq \frac{1}{4} \left( \left| \frac{H \delta - a_p}{2} \right| + \left| \frac{H \delta - a_q}{2} \right| \right) \]

\[ \geq \left| \frac{a_p - a_q}{4} \right| \geq \left| \frac{c_I^{(p)}(f) - c_I^{(q)}(f)}{4} \right|, \]

which proves the last statement of the lemma.
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