Towards Scalable Synthesis of Stochastic Control Systems

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Abstract Formal synthesis approaches over stochastic systems have received significant attention in the past few years, in view of their ability to provide provably correct controllers for complex logical specifications in an automated fashion. Examples of complex specifications include properties expressed as formulae in linear temporal logic (LTL) or as automata on infinite strings. A general methodology to synthesize controllers for such properties resorts to symbolic models of the given stochastic systems. Symbolic models are finite abstractions of the given concrete systems with the property that a controller designed on the abstraction can be refined (or implemented) into a controller on the original system. Although the recent development of techniques for the construction of symbolic models has been quite encouraging, the general goal of formal synthesis over stochastic control systems is by no means solved. A fundamental issue with the existing techniques is the known “curse of dimensionality,” which is due to the need to discretize state and input sets. Such discretization generally results in an exponential complexity over the number of state and input variables in the concrete system. In this work we propose a novel abstraction technique for incrementally stable stochastic control systems, which does
not require state-space discretization but only input set discretization, and that can be potentially more efficient (and thus scalable) than existing approaches. We elucidate the effectiveness of the proposed approach by synthesizing a schedule for the coordination of two traffic lights under some safety and fairness requirements for a road traffic model. Further we argue that this 5-dimensional linear stochastic control system cannot be studied with existing approaches based on state-space discretization due to the very large number of generated discrete states.

**Keywords** Stochastic control systems · Formal controller synthesis · Finite abstractions · Approximate bisimulation

1 Introduction

In the last decade many techniques have been developed providing controllers for control systems (both deterministic and, more recently, stochastic) in a formal and automated fashion against some complex logical specifications. Examples of such specifications include properties expressed as formulae in linear temporal logic (LTL) or as automata on infinite strings [3], and as such they are not tractable by classical techniques for control systems. A general scheme for providing such controllers is by leveraging symbolic models of original concrete systems. Symbolic models are discrete abstractions of the original systems in which each symbol represents an aggregate of continuous variables. When such symbolic models exist for the concrete systems, one can leverage the algorithmic machinery for automated synthesis of discrete models [1,14] to automatically synthesize discrete controllers which can be refined to hybrid controllers for the original systems.

The construction of symbolic models for continuous-time deterministic systems has been thoroughly investigated in the past few years. This includes results on the construction of approximately bisimilar symbolic models for incrementally stable control systems [15,18], incrementally stable switched systems [7], and control systems with disturbances [19], non-uniform abstractions of nonlinear systems over a finite-time horizon [26], as well as the construction of sound abstractions based on the convexity of reachable sets [20], feedback refinement relations [21], robustness margins [13], and approximate alternating simulation relation [32]. Recently, there have been some results on the construction of symbolic models for continuous-time stochastic systems, including the construction of finite Markov decision process of linear stochastic control system, without providing a quantitative relationship between abstract and concrete model [12], approximately bisimilar symbolic models for incrementally stable stochastic control systems [31], stochastic switched systems [29], and randomly switched stochastic systems [28], as well as abstractions for unstable stochastic control systems [30].

Note that all the techniques provided in [15,18,7,19,26,20,21,13,32,12,31,28,30] are fundamentally based on the discretization of continuous states. Therefore, they suffer severely from the curse of dimensionality due to gridding those sets, which is especially irritating for models with high-dimensional state spaces. In this work we propose a novel approach for the construction of approximately bisimilar symbolic models for incrementally stable stochastic control systems not requiring any state set.
discretization but only input set discretization. Therefore, it can be potentially more efficient than the proposed approaches in [31] when dealing with higher dimensional stochastic control systems. In particular, we provide a theoretical comparison with the approach in [31] and a simple criterion that helps choosing the most suitable among two approaches (in terms of the sizes of the symbolic models) for a given stochastic control system.

Another advantage of the technique proposed here is that it allows us to construct symbolic models with probabilistic output values, resulting in less conservative symbolic abstractions than those proposed in [31,28,30] that allow for deterministic output values exclusively. We then explain how the proposed symbolic models with probabilistic output values can be used for synthesizing hybrid controllers enforcing logic specifications. The proposed approaches in [29] also provide symbolic models with probabilistic output values and without any state set discretization. However, the results in [29] are for stochastic switched systems rather than stochastic control systems as in this work and they do not provide any intuition behind the control synthesis over symbolic models with probabilistic output values. The effectiveness of the proposed results is illustrated by synthesizing a schedule for the coordination of two traffic lights under some safety and fairness requirements for a model of road traffic which is a 5-dimensional linear stochastic control system. We also show that this example is not amenable to be dealt with the approaches proposed in [31]. Although the main proposed results in this work are for incrementally stable stochastic control systems, the similar results for incrementally stable deterministic control systems can be recovered in the same framework by simply setting the diffusion term to zero.

Alongside the relationship with and extension of [29,31], this paper provides a detailed and extended elaboration of the results first announced in [34], including the proofs of the main results, a detailed discussion on how to deal with probabilistic output values and a generalization of the corresponding result with no requirement on compactness, and finally discussing a new case study on road traffic control.

Similar to this contribution, the results in [23,25] also use sequences of symbols as symbolic states to provide finite abstractions of concrete systems. However, the results in [23,25] apply to discrete-time deterministic control systems, rather than continuous-time stochastic control systems. Furthermore, the results in [25] do not provide a constructive approach for the synthesis of finite abstractions, and the relationships established between abstractions and concrete systems do not necessarily preserve LTL properties. The results in [23] provide asynchronous ℓ-complete finite abstractions which are generally incomparable to the approximate bisimilar finite abstractions proposed in this work.

2 Stochastic Control Systems

2.1 Notation

The identity map over a set $B$ is denoted by $1_B$. The symbols $\mathbb{N}$, $\mathbb{N}_0$, $\mathbb{Z}$, $\mathbb{R}$, $\mathbb{R}^+$, and $\mathbb{R}_0^+$ denote the set of natural, nonnegative integer, integer, real, positive, and nonnegative real numbers, respectively. The symbols $I_n$, $0_n$, and $0_{n \times m}$ denote the identity matrix,
the zero vector, and the zero matrix in \( \mathbb{R}^{n \times n} \), \( \mathbb{R}^n \), and \( \mathbb{R}^{n \times m} \), respectively. Given a vector \( x \in \mathbb{R}^n \), we denote by \( x_i \) the \( i \)-th element of \( x \), and by \( \| x \| \) the infinity norm of \( x \). Given a matrix \( P = \{ p_{ij} \} \in \mathbb{R}^{n \times n} \), we denote by \( \text{Tr}(P) = \sum_{i=1}^{n} p_{ii} \) the trace of \( P \).

Given a symmetric matrix \( A \), we denote by \( \lambda_{\text{min}}(A) \) and \( \lambda_{\text{max}}(A) \) the minimum and maximum eigenvalues of \( A \), respectively. The diagonal set \( \Delta \subset \mathbb{R}^n \times \mathbb{R}^n \) is defined as:

\[
\Delta = \{(x, x) \mid x \in \mathbb{R}^n \}.
\]

The closed ball centered at \( u \in \mathbb{R}^m \) with radius \( \lambda \) is defined by \( B_\lambda(u) = \{ v \in \mathbb{R}^m \mid \| u - v \| \leq \lambda \} \). A set \( B \subset \mathbb{R}^m \) is called a box if \( B = \prod_{i=1}^{m} [c_i, d_i] \), where \( c_i, d_i \in \mathbb{R} \) with \( c_i < d_i \) for each \( i \in \{1, \ldots, m\} \). The span of a box \( B \) is defined as \( \text{span}(B) = \min \{ |d_i - c_i| \mid i = 1, \ldots, m \} \). For a box \( B \subset \mathbb{R}^m \) and \( \mu \leq \text{span}(B) \), define the \( \mu \)-approximation \( [B]_\mu = [\mathbb{R}^m]_\mu \cap B \), where \( [\mathbb{R}^m]_\mu = \{ a \in \mathbb{R}^m \mid a_i = k_i \mu, k_i \in \mathbb{Z}, i = 1, \ldots, m \} \). Remark that \( [B]_\mu \neq \emptyset \) for any \( \mu \leq \text{span}(B) \). Geometrically, for any \( \mu \in \mathbb{R}^+ \) with \( \mu \leq \text{span}(B) \) and \( \lambda \geq \mu \), the collection of sets \( \{ B_\lambda(p) \} \in [\mathbb{R}]_\mu \) is a finite covering of \( B \), i.e. \( B \subset \bigcup_{p \in [\mathbb{R}]_\mu} B_\lambda(p) \). We extend the notions of \( \text{span} \) and \( \text{approximation} \) to finite unions of boxes as follows. Let \( A = \bigcup_{j=1}^{m} A_j \), where each \( A_j \) is a box. Define \( \text{span}(A) = \min \{ \text{span}(A_j) \mid j = 1, \ldots, M \} \), and for any \( \mu \leq \text{span}(A) \), define \( \text{span}(A) = \bigcup_{j=1}^{m} [A_j]_\mu \).

Given a measurable function \( f : \mathbb{R}^n_+ \to \mathbb{R}^n \), the (essential) supremum of \( f \) is denoted by \( \| f \|_\infty = (\text{ess}\sup) \{ \| f(t) \|, t \geq 0 \} \). A continuous function \( \gamma : \mathbb{R}^n_+ \to \mathbb{R}^n_+ \) is said to belong to class \( \mathcal{X} \) if it is strictly increasing and \( \gamma(0) = 0 \); \( \gamma \) is said to belong to class \( \mathcal{X}_\infty \) if \( \gamma \in \mathcal{X} \) and \( \gamma(r) \to \infty \) as \( r \to \infty \). A continuous function \( \beta : \mathbb{R}^n_+ \times \mathbb{R}^n_+ \to \mathbb{R}^n_+ \) is said to belong to class \( \mathcal{X} \mathcal{X} \) if, for each fixed \( s \), the map \( \beta(r, s) \) belongs to class \( \mathcal{X} \) with respect to \( r \) and, for each fixed nonzero \( r \), the map \( \beta(r, s) \) is decreasing with respect to \( s \) and \( \beta(r, s) \to 0 \) as \( s \to \infty \). We identify a relation \( R \subseteq A \times B \) with the map \( R : A \to \mathbb{Z}^d \) defined by \( b \in R(a) \) iff \( (a, b) \in R \). Given a relation \( R \subseteq A \times B \), \( R^{-1} \) denotes the inverse relation defined by \( R^{-1} = \{(b, a) \in B \times A \mid (a, b) \in R \} \). Given a finite sequence \( S \), we denote by \( \sigma := (S)^\omega \) the infinite sequence generated by repeating \( S \) infinitely, i.e. \( \sigma := \underline{SSSSSSS} \ldots \).

2.2 Stochastic control systems

Let \( (\Omega, \mathcal{F}, \mathbb{P}) \) be a probability space equipped with a filtration \( \mathbb{F} = (\mathcal{F}_t)_{t \geq 0} \) satisfying the usual conditions of completeness and right continuity [10, p. 48]. Let \( (W_t)_{t \geq 0} \) be a \( p \)-dimensional \( \mathbb{F} \)-adapted Brownian motion.

**Definition 1** A stochastic control system \( \Sigma \) is a tuple \( \Sigma = (\mathbb{R}^n, \mathcal{U}, \mathcal{V}, f, \sigma) \), where \( \mathbb{R}^n \) is the state space, \( \mathcal{U} \subset \mathbb{R}^m \) is a bounded input set, and

- \( \mathcal{V} \) is a subset of the set of all measurable functions from \( \mathbb{R}^n_+ \) to \( \mathcal{U} \);
- \( f : \mathbb{R}^n_+ \times \mathcal{U} \to \mathbb{R}^n \) satisfies the following Lipschitz assumption: there exist constants \( L_x, L_u \in \mathbb{R}^+ \) such that \( \| f(x, u) - f(x', u') \| \leq L_x \| x - x' \| + L_u \| u - u' \| \) for all \( x, x' \in \mathbb{R}^n \) and all \( u, u' \in \mathcal{U} \);
- \( \sigma : \mathbb{R}^n \to \mathbb{R}^{n \times p} \) satisfies the following Lipschitz assumption: there exists a constant \( Z \in \mathbb{R}^+ \) such that \( \| \sigma(x) - \sigma(x') \| \leq Z \| x - x' \| \) for all \( x, x' \in \mathbb{R}^n \). \( \square \)
A continuous-time stochastic process $\xi : \Omega \times \mathbb{R}_0^+ \to \mathbb{R}^n$ is said to be a solution process of $\Sigma$ if there exists $\upsilon \in \mathcal{U}$ satisfying the following stochastic differential equation (SDE) $\mathbb{P}$-almost surely ($\mathbb{P}$-a.s.):

$$d\xi = f(\xi, \upsilon)\,dt + \sigma(\xi)\,dW_t,$$

where $f$ is known as the drift and $\sigma$ as the diffusion. We also denote by $\xi_{\sigma_0}(t)$ the value of the solution process at time $t \in \mathbb{R}_0^+$ under the input curve $\upsilon$ from initial condition $\xi_{\sigma_0}(0) = a$ $\mathbb{P}$-a.s., where $a$ is a $\mathcal{F}_0$-measurable random variable. Let us emphasize that the solution process is unambiguously determined, since the assumptions on $f$ and $\sigma$ ensure its existence and uniqueness [16, Theorem 5.2.1, p. 68].

**Example 1** For a linear stochastic control system $\Sigma = (\mathbb{R}^n, U, \mathcal{W}, f, \sigma)$, one has $f(x, u) := Ax + Bu$, for some matrices $A \in \mathbb{R}^{n \times n}$ and $B \in \mathbb{R}^{n \times m}$; and $\sigma(x) := [\sigma_1, \ldots, \sigma_p]$, for some matrices $\sigma_i \in \mathbb{R}^{n \times n}$.

### 3 Incremental Stability

We recall a stability notion for stochastic control systems, introduced in [31], on which the main results presented in this work rely.

**Definition 2** A stochastic control system $\Sigma$ is incrementally input-to-state stable in the $q$th moment ($\delta$-ISS-M$_q$), where $q \geq 1$, if there exist a function $\beta \in \mathcal{K} \mathcal{L}_q$ and a function $\gamma \in \mathcal{K}_\infty$ such that for any $t \in \mathbb{R}_0^+$, any $\mathbb{R}^n$-valued $\mathcal{F}_0$-measurable random variables $a$ and $a'$, and any $\upsilon$, $\upsilon' \in \mathcal{W}$, the following condition is satisfied:

$$E[\|\xi_{\sigma_0}(t) - \xi_{\sigma_1}(t)\|_q^q] \leq \beta(E[\|a - a'\|_q^q], t) + \gamma(\|\upsilon - \upsilon'\|_\infty).$$

It can be easily verified that a $\delta$-ISS-M$_q$ stochastic control system $\Sigma$ is $\delta$-ISS [2] in the absence of any noise as in the following:

$$\|\xi_{\sigma_0}(t) - \xi_{\sigma_1}(t)\| \leq \beta(E[\|a - a'\|_q], t) + \gamma(\|\upsilon - \upsilon'\|_\infty),$$

for $a, a' \in \mathbb{R}^n$, some $\beta \in \mathcal{K} \mathcal{L}_q$, and some $\gamma \in \mathcal{K}_\infty$.

**Remark 1** For linear stochastic control systems $\Sigma$, as in Example 1, $\delta$-ISS-M$_2$ is the same as asymptotic stability in the mean square sense (cf. [33, Lemma 5.1]). However, for nonlinear stochastic control systems, $\delta$-ISS-M$_q$ is a much stronger property than the usual stability in the $q$th moment [9].

Similar to the characterization of $\delta$-ISS in terms of the existence of so-called $\delta$-ISS Lyapunov functions in [2], one can describe $\delta$-ISS-M$_q$ in terms of the existence of so-called $\delta$-ISS-M$_q$ Lyapunov functions, as shown in [31] and defined next.

**Definition 3** Consider a stochastic control system $\Sigma$ and a continuous function $V : \mathbb{R}_0^+ \times \mathbb{R}^n \to \mathbb{R}_0^+$ that is twice continuously differentiable on $\{\mathbb{R}_0^+ \times \mathbb{R}^n\} \setminus \Delta$. The function $V$ is called a $\delta$-ISS-M$_q$ Lyapunov function for $\Sigma$, where $q \geq 1$, if there exist functions $\varphi, \tau, \rho$ in $\mathcal{K}_\infty$, and a constant $\kappa \in \mathbb{R}_0^+$, such that...
(i) $\varphi$ (resp. $\overline{\varphi}$) is a convex (resp. concave) function; 
(ii) for any $x, x' \in \mathbb{R}^n$, $\varphi(||x - x'||^q) \leq V(x, x') \leq \overline{\varphi}(||x - x'||^q)$; 
(iii) for any $x, x' \in \mathbb{R}^n$, $x \neq x'$, and for any $u, u' \in U$, 
\[ L_{u, u'} V(x, x') = \left[ \frac{\partial}{\partial x} V \frac{\partial}{\partial x'} V \right] \left[ \begin{array}{c} f(x, u) \\ f(x', u') \end{array} \right] + \frac{1}{2} \text{Tr} \left( \left[ \sigma(x) \sigma(x') \right] \left[ \sigma^T(x) \sigma^T(x') \right] \left[ \frac{\partial}{\partial x} V \frac{\partial}{\partial x'} V \right] \right) \leq -\kappa V(x, x') + \rho (||u - u'||), \]

where $L_{u, u'}$ is the infinitesimal generator associated to the process $V(\xi, \xi')$ and $\xi$ and $\xi'$ are solution processes of the SDE (1) [16, Section 7.3]. The symbols $\frac{\partial}{\partial x}$ and $\frac{\partial}{\partial x'}$ denote first- and second-order partial derivatives with respect to $x$ and $(x, x')$, respectively. 

Condition (ii) in the above definition implies that the growth rate of functions $\overline{\varphi}$ and $\varphi$ is linear. Remark that this condition does not limit the behavior of $\varphi$ and $\overline{\varphi}$ to linear ones on a compact subset of $\mathbb{R}^n$. Observe that condition (i) is not required in the context of deterministic control systems for the corresponding $\delta$-ISS Lyapunov functions [2]. The following theorem, borrowed from [31], describes $\delta$-ISS-$M_q$ in terms of the existence of $\delta$-ISS-$M_q$ Lyapunov functions.

**Theorem 1** A stochastic control system $\Sigma$ is $\delta$-ISS-$M_q$ if it admits a $\delta$-ISS-$M_q$ Lyapunov function. 

One can resort to available optimization software tools, such as SOSTOOLS [17], to search for appropriate $\delta$-ISS-$M_q$ Lyapunov functions for systems $\Sigma$ of polynomial type (polynomial drift and diffusion). We refer the interested readers to the results in [31] for the discussion of special instances where these functions can be easily computed.

As an example, for linear stochastic control systems $\Sigma$, as in Example 1, one can search for $\delta$-ISS-$M_q$ Lyapunov functions of the form $V(x, x') := \frac{1}{q}((x - x')^T P (x - x'))$, for some positive definite $P \in \mathbb{R}^{n \times n}$ and $q \in \{1, 2\}$, by solving the following linear matrix inequality (LMI):

\[ A^T P + PA + \sum_{i=1}^{n} \sigma_i^T P \sigma_i \preceq -\hat{\kappa}P, \quad (4) \]

for some positive constant $\hat{\kappa}$. We refer the interested readers to [31, Corollary 3.5] for more details.

### 3.1 Noisy and noise-free trajectories

In order to introduce the symbolic models with deterministic output values in Subsection 5.2 (Theorems 2 and 3) for a stochastic control system, we need the following technical result, borrowed from [31]. The following result provides an upper bound on the distance (in the $q$th moment) between the solution process of
Σ and the solution of a derived deterministic control system \( \bar{\Sigma} \) obtained by disregarding the diffusion term \( \sigma \) in \( \Sigma \). From now on, we denote by \( \bar{\xi}(\xi, \upsilon) \) the solution of \( \bar{\Sigma} = (\mathbb{R}^n, \mathcal{U}, \mathcal{F}, 0_{n \times p}) \), starting from the deterministic initial condition \( \xi = \bar{\xi}(\xi, 0) \) and under the input curve \( \upsilon \), which satisfies the ordinary differential equation (ODE) \( \dot{\xi}(\xi, \upsilon) = f(\xi, \upsilon) \).

**Lemma 1** Consider a stochastic control system \( \Sigma \) such that \( f(0_n, 0_m) = 0_n \) and \( \sigma(0_n) = 0_{n \times p} \). Suppose that \( q \geq 2 \) and that there exists a \( \delta \)-ISS-M Lyapunov function \( V \) for \( \Sigma \) such that its Hessian is a positive semidefinite matrix in \( \mathbb{R}^{2n \times 2n} \) and \( \partial_x V(x, x') \leq P \), for any \( x, x' \in \mathbb{R}^n \), and some positive semidefinite matrix \( P \in \mathbb{R}^{n \times n} \).

Then for any \( x \in \mathbb{R}^n \) and any \( \upsilon \in \mathcal{U} \), we have

\[
\mathbb{E} \left[ \left\| \xi(\xi, \upsilon(t)) - \bar{\xi}(\xi, \upsilon(t)) \right\|^q \right] \leq h_x(t),
\]

where

\[
h_x(t) = \frac{\alpha}{2} \left( \frac{1}{2} \sqrt{\|P\|} \right)^2 \min\{n, p\} Z^2 e^{-\kappa t} \int_0^t \left( \beta \left( \|x\|, s \right) + \gamma \left( \sup_{u \in \mathcal{U}} \{\|u\|\} \right) \right)^{\frac{1}{q}} ds,
\]

and \( Z \) is the Lipschitz constant, introduced in Definition 1, and \( \beta \) is the \( \mathcal{KL} \) function appearing in (2).

It can be readily seen that the nonnegative-valued function \( h_x \) tends to zero as \( t \to 0, t \to +\infty \), or as \( Z \to 0 \), and is identically zero if the diffusion term is identically zero (i.e. \( Z = 0 \), which is the case for \( \bar{\Sigma} \)). The interested readers are referred to [31] providing (possibly less conservative) results in line with that of Lemma 1 for (linear) stochastic control systems \( \Sigma \) admitting a specific type of \( \delta \)-ISS-M Lyapunov functions.

### 4 Systems and Approximate Equivalence Relations

#### 4.1 Systems

We employ the abstract and general notion of “system,” as introduced in [24], to describe both stochastic control systems and their symbolic models.

**Definition 4** A system \( S \) is a tuple \( S = (X, X_0, U, \rightarrow, Y, H) \), where \( X \) is a (possibly infinite) set of states, \( X_0 \subseteq X \) is a (possibly infinite) set of initial states, \( U \) is a (possibly infinite) set of inputs, \( \rightarrow \subseteq X \times U \times X \) is a transition relation, \( Y \) is a set of outputs, and \( H : X \to Y \) is an output map. A transition \( (x, u, x') \in \rightarrow \) is also denoted by \( x \xrightarrow{u} x' \). For a transition \( x \xrightarrow{u} x' \), state \( x' \) is called a \( u \)-successor, or simply a successor, of state \( x \). We denote by \( Post_u(x) \) the set of all \( u \)-successors of a state \( x \). A system \( S \) is said to be

- *metric*, if the output set \( Y \) is equipped with a metric \( d : Y \times Y \to \mathbb{R}_0^+ \);

\( ^1 \) Here, we have abused notation by identifying \( 0_{n \times p} \) with the map \( \sigma : x \to 0_{n \times p} \forall x \in \mathbb{R}^n \).
– finite (or symbolic), if $X$ and $U$ are finite sets;
– deterministic, if for any state $x \in X$ and any input $u \in U$, $|\text{Post}_u(x)| \leq 1$. □

For technical reasons, we assume that for any $x \in X$, there exists some $u$-successor of $x$, for some $u \in U$ — note that this is always the case for the considered systems in this paper.

Given a system $S = (X, X_0, U, \longrightarrow , Y, H)$ and any initial state $x_0 \in X_0$, a finite state run generated from $x_0$ is a finite sequence of transitions:

$$x_0 \xrightarrow{u_0} x_1 \xrightarrow{u_1} \cdots x_{n-1} \xrightarrow{u_{n-1}} x_n,$$

such that $x_i \xrightarrow{u_i} x_{i+1}$ for all $0 \leq i < n$. A finite state run can be directly extended to an infinite state run as well. A finite output run is a sequence $\{y_0, y_1, \ldots, y_n\}$ such that there exists a finite state run of the form (6) with $y_i = H(x_i)$, for $i = 0, \ldots, n$. A finite output run can also be directly extended to an infinite output run as well.

### 4.2 Relations among systems

We recall the notion of approximate (bi)simulation relation, introduced in [6], which is crucial when analyzing or synthesizing controllers for deterministic systems.

**Definition 5** Consider two metric systems $S_a = (X_a, X_{0a}, U_a, \longrightarrow^a , Y_a, H_a)$ and $S_b = (X_b, X_{0b}, U_b, \longrightarrow^b , Y_b, H_b)$ with the same output sets $Y_a = Y_b$ and metric $d$. For $\varepsilon \in \mathbb{R}_0^+$, a relation $R \subseteq X_a \times X_b$ is said to be an $\varepsilon$-approximate simulation relation from $S_a$ to $S_b$ if, for all $(x_a, x_b) \in R$, the following two conditions are satisfied:

(i) $d(H_a(x_a), H_b(x_b)) \leq \varepsilon$;
(ii) for any $x_a \xrightarrow{a} x'_a \in S_a$, there exists $x_b \xrightarrow{b} x'_b \in S_b$ such that $(x'_a, x'_b) \in R$.

A relation $R \subseteq X_a \times X_b$ is called an $\varepsilon$-approximate bisimulation relation between $S_a$ and $S_b$ if $R$ is an $\varepsilon$-approximate simulation relation from $S_a$ to $S_b$ and $R^{-1}$ is an $\varepsilon$-approximate simulation relation from $S_b$ to $S_a$.

System $S_a$ is $\varepsilon$-approximately simulated by $S_b$, or $S_b$ $\varepsilon$-approximately simulates $S_a$, denoted by $S_a \preceq^\varepsilon \prec S_b$, if there exists an $\varepsilon$-approximate simulation relation $R$ from $S_a$ to $S_b$ such that:

– for every $x_{a0} \in X_{a0}$, there exists $x_{b0} \in X_{b0}$ such that $(x_{a0}, x_{b0}) \in R$.

System $S_a$ is $\varepsilon$-approximately bisimilar to $S_b$, denoted by $S_a \cong^\varepsilon \prec S_b$, if there exists an $\varepsilon$-approximate bisimulation relation $R$ between $S_a$ and $S_b$ such that:

– for every $x_{a0} \in X_{a0}$, there exists $x_{b0} \in X_{b0}$ such that $(x_{a0}, x_{b0}) \in R$.
– for every $x_{a0} \in X_{a0}$, there exists $x_{b0} \in X_{b0}$ such that $(x_{a0}, x_{b0}) \in R$. □
5 Symbolic Models for Stochastic Control Systems

5.1 Describing stochastic control systems as metric systems

In order to show the main results of the paper, we use the notion of system introduced in Definition 4 to abstractly represent a stochastic control system. More precisely, given a stochastic control system \( \Sigma \), we define an associated metric system \( S(\Sigma) = (X, X_0, U, Y, H) \), where:

- \( X \) is the set of all \( \mathbb{R}^n \)-valued random variables defined on the probability space \( (\Omega, \mathcal{F}, \mathbb{P}) \);
- \( X_0 \) is a subset of the set of \( \mathbb{R}^n \)-valued random variables that are measurable over \( \mathcal{F}_0 \);
- \( U = \mathcal{U} \);
- \( x \) and \( x' \) are measurable in \( \mathcal{F}_t \) and \( \mathcal{F}_{t+\tau} \), respectively, for some \( t \in \mathbb{R}^+ \) and \( \tau \in \mathbb{R}^+ \), and there exists a solution process \( \xi : \Omega \times \mathbb{R}^+_0 \to \mathbb{R}^n \) of \( \Sigma \) satisfying \( \xi(t) = x \) and \( \xi_{x'}(\tau) = x' \) \( \mathbb{P} \)-a.s.;
- \( Y = X \);
- \( H = 1_X \).

We assume that the output set \( Y \) is equipped with the \( q \)th moment metric \( d(y,y') = (E[\|y-y'\|^q])^{\frac{1}{q}} \), for any \( y,y' \in Y \) and some \( q \geq 1 \). One can readily observe that the set of states and inputs of \( S(\Sigma) \) are uncountable and that \( S(\Sigma) \) is a deterministic system in the sense of Definition 4, since (cf. Subsection 2.2) the solution process of \( \Sigma \) is uniquely determined. For the case of deterministic control system \( \Sigma \), one obtains \( S(\Sigma) = (X, X_0, U, Y, H) \), where \( X = \mathbb{R}^n \), \( X_0 \) is a subset of \( \mathbb{R}^n \), \( U = \mathcal{U} \), \( x \) if \( x' = \xi_{x'}(\tau) \) for some \( \tau \in \mathbb{R}^+ \), \( Y = X \), \( H = 1_X \), and the metric on the output set reduces to the natural infinity one: \( d(y,y') = \|y-y'\| \), for any \( y,y' \in Y \).

Since the concrete system \( S(\Sigma) \) is uncountably infinite, it does not allow for a straightforward symbolic controller synthesis with the discrete techniques in the literature [1,14]. We are thus interested in finding a finite abstract system that is (bi)similar to the concrete one \( S(\Sigma) \). In order to discuss approximate (bi)simulation relations between two metric systems, they have to share the output space (cf. Definition 5). System \( S(\Sigma) \) inherits a classical trace-based semantics [3] (cf. definition of output run after (6)), however the outputs of \( S(\Sigma) \) (and necessarily those of any approximately (bi)similar one) are random variables. This fact is especially important due to the metric \( d \) that the output set is endowed with: for any deterministic point one can always find a non-degenerate random variable that is as close as desired to the original point in the metric \( d \).

To further elaborate the discussion in the previous paragraph, let us consider the following example. Let \( A \subset \mathbb{R}^n \) be a set (of deterministic points). Consider a safety problem, formulated as the satisfaction of the LTL formula\(^2\) \( \square \phi_A \), where \( \phi_A \) is a label (or proposition) characterising the set \( A \). Suppose that over the abstract system we are able to synthesize a control strategy that makes an output run of the abstraction satisfy

\(^2\) We refer the interested readers to [3, Subsection 5.1.2] for the formal trace-based semantic of LTL formulae.
\[ \square \varphi_A , \] which means that the output run of the abstraction is always inside the set \( A \).

Although the run would in general be consisting of random variables \( y : \Omega \to \mathbb{R}^n \), the fact that \( y \in A \) means that, with a slight abuse of notation, \( y : \bar{\Omega} \to \{ y \} \) where \( \bar{\Omega} \subset \Omega \) and \( |\bar{\Omega}| = 1 \). Hence, \( y \) has a Dirac probability distribution centered at \( y \), that is \( y \in Y \) is a degenerate random variable that can be identified with a point in \( A \subset \mathbb{R}^n \subset Y \).

Since any deterministic point can be regarded as a random variable with a Dirac probability distribution centered at that point, \( \mathbb{R}^n \) can be embedded in \( Y \), which we denote as \( \mathbb{R}^n \subset Y \) with a slight abuse of notation. As a result, satisfying \( \square \varphi_A \) precisely means that the output run of the abstraction indeed stays in the set \( A \subset \mathbb{R}^n \) forever.

On the other hand, suppose that the original system is \( \varepsilon \)-approximate bisimilar to \( \varphi_A \).

Although the run would in general be consisting of random variables \( y : \Omega \to \mathbb{R}^n \), the fact that \( y \in A \) means that, with a slight abuse of notation, \( y : \bar{\Omega} \to \{ y \} \) where \( \bar{\Omega} \subset \Omega \) and \( |\bar{\Omega}| = 1 \). Hence, \( y \) has a Dirac probability distribution centered at \( y \), that is \( y \in Y \) is a degenerate random variable that can be identified with a point in \( A \subset \mathbb{R}^n \subset Y \).

We are now able to provide two versions of finite abstractions: one whose outputs are always deterministic points, that is any output \( y \) of the run of the original system is within \( \varepsilon \) \( d \)-distance from the set \( A : d(y, A) = \inf_{a \in A} d(y, a) \leq \varepsilon \). Although the original set \( A \subset Y \) is a subset of \( \mathbb{R}^n \subset Y \), its \( \varepsilon \)-inflation \( A_\varepsilon = \{ y \in Y : d(y, A) \leq \varepsilon \} \) is not a subset of \( \mathbb{R}^n \) anymore and hence contains non-degenerate random variables. In particular, \( A_\varepsilon \neq \{ y \in \mathbb{R}^n : \inf_{a \in A} \| y - a \| \leq \varepsilon \} \) and is in fact bigger than the latter set of deterministic points.

As a result, although satisfying \( \square \varphi_A \) does not necessarily mean that a trajectory of \( \Sigma \) always stays within some deterministic set, it means that the associated random variables always belong to \( A_\varepsilon \) and, hence, are close to the deterministic set \( A \) with respect to the \( q \)th moment metric.

5.2 Main results

This subsection contains the main contributions of the paper. We show that for any \( \delta \)-ISS-\( M_\gamma \) (resp. \( \delta \)-ISS) stochastic control system \( \Sigma \) (resp. deterministic control system \( \Sigma \)), and for any precision level \( \varepsilon \in \mathbb{R}^+ \), we can construct a finite system that is \( \varepsilon \)-approximate bisimilar to \( \Sigma \) (resp. \( \Sigma \)) without any state set discretization. The results in this subsection rely on additional assumptions on the model \( \Sigma \) that are described next. We restrict our attention to stochastic control systems \( \Sigma \) with input sets \( U \) that are assumed to be finite unions of boxes (cf. Subsection 2.1). We further restrict our attention to sampled-data stochastic control systems, where input curves belong to set \( \mathcal{W}_\tau \), which contains exclusively curves that are constant over intervals of length \( \tau \in \mathbb{R}^+ \), i.e.

\[
\mathcal{W}_\tau = \left\{ v \in \mathcal{W} \mid v(t) = v((k-1)\tau), t \in [(k-1)\tau, k\tau], k \in \mathbb{N} \right\}.
\]

Let us denote by \( S_\varepsilon(\Sigma) \) a sub-system of \( S(\Sigma) \) obtained by selecting those transitions of \( S(\Sigma) \) corresponding to solution processes of duration \( \tau \) and to control inputs in \( \mathcal{W}_\tau \).
This can be seen as the sampled version of $\Sigma$. More precisely, given a stochastic control system $\Sigma$ and the corresponding metric system $S(\Sigma)$, we define a new associated metric system

$$S_\varepsilon(\Sigma) = \left( X_\varepsilon, X_{\tau 0}, U_\varepsilon, \right.$$

$$\left. \xrightarrow{\varepsilon} Y_\varepsilon, H_\varepsilon \right),$$

where $X_\varepsilon = X, X_{\tau 0} = X_0, U_\varepsilon = U_\tau, Y_\varepsilon = Y, H_\varepsilon = H$, and $x_t \xrightarrow{\varepsilon} x'_t \in X$ if $x_t$ and $x'_t$ are measurable, respectively, in $\mathcal{F}_t$ and $\mathcal{F}_{(k+1)\tau}$ for some $k \in \mathbb{N}_0$, and there exists a solution process $\xi : \Omega \times \mathbb{R}^+_0 \to \mathbb{R}^n$ of $\Sigma$ satisfying $\xi(k\tau) = x_t$ and $\xi_{x_t u_\varepsilon}(\tau) = x'_t \mathbb{P}$-a.s.

Similarly, one can define $S_\varepsilon^*(\Sigma)$ as the time discretization of $\Sigma$. Observe that a finite state run $x_0 \xrightarrow{u_0} x_1 \xrightarrow{u_1} \cdots x_N$ of $S(\Sigma)$, where $u_{i-1} \in U_\tau$ and $x_i = \xi_{u_{i-1} u_i}(\tau)$, captures the solution process of $\Sigma$ at times $t = 0, \tau, \ldots, N\tau$, started from the initial condition $x_0$ and resulting from a control input $u$ obtained by the concatenation of the input curves $u_{i-1}$ (i.e. $u(t) = u_{i-1}(t)$ for any $t \in [(i-1)\tau, i\tau]$), for $i = 1, \ldots, N$.

Now, we have all the ingredients to introduce two fully symbolic systems for the concrete model $\Sigma$. Consider a stochastic control system $\Sigma$ and a tuple $q = (\tau, \mu, N, x_0)$ of parameters, where $\tau$ is the sampling time, $\mu$ is the input set quantization, $N \in \mathbb{N}$ is a temporal horizon, and $x_0 \in \mathbb{R}^n$ is a source state. Given $\Sigma$ and $q$, let us introduce the following two symbolic systems:

$$S_q(\Sigma) = \left( X_q, X_{q0}, U_q, \xrightarrow{q} Y_q, H_q \right),$$

$$\overline{S}_q(\Sigma) = \left( X_q, X_{q0}, U_q, \xrightarrow{q} Y_q, \overline{H}_q \right),$$

where

- $X_q := \{ (u_1, \ldots, u_N) \in \lceil 1 \rceil \times \cdots \times \lceil 1 \rceil \}$;
- $X_{q0} = X_0$;
- $U_q = \lceil 1 \rceil$;
- $x_q \xrightarrow{u_q} x'_q$, where $x_q = (u_1, u_2, \ldots, u_N)$, if and only if $x'_q = (u_2, \ldots, u_N, u_0)$;
- $Y_q$ is the set of all $\mathbb{R}^n$-valued random variables defined on the probability space $(\Omega, \mathcal{F}, \mathbb{P})$;
- $H_q(x_q) = \xi_{x_q}(N\tau)$, $\overline{H}_q(x_q) = \overline{\xi}_{x_q}(N\tau)$.

Note that the transition relation in $S_q(\Sigma)$ admits a compact representation in the form of a shift operator. We have abused notation by identifying $u_q \in \lceil 1 \rceil$ with the constant input curve with domain $[0, \tau]$ and value $u_0$, and by identifying $x_0 \in \lceil 1 \rceil^{\times N}$ with the concatenation of $N$ control inputs $u_i \in \lceil 1 \rceil$ (i.e. $x_q(t) = u_i$ for any $t \in [(i-1)\tau, i\tau]$) for $i = 1, \ldots, N$. Notice that the proposed abstraction $S_q(\Sigma)$ (resp. $\overline{S}_q(\Sigma)$) is a deterministic system in the sense of Definition 4. We point out that $H_q$ and $\overline{H}_q$ are mappings from a deterministic point $x_0$ to the random variable $\xi_{x_0}(N\tau)$ and to the one with a Dirac probability distribution centered at $\overline{\xi}_{x_0}(N\tau)$, respectively. Finally,
in the case of a deterministic control system $\Sigma$, one obtains the symbolic system $\bar{S}_q(\Sigma) = (X_q, X_{\bar{q}}, U_q, Q_q(\Sigma), Y_q, \bar{Y}_q)$, where $X_q, X_{\bar{q}}, U_q, Q_q(\Sigma), Y_q, \bar{Y}_q$ are the same as before, but the output set reduces to $Y_q = \mathbb{R}^n$.

The main idea behind the definitions of symbolic models $S_q(\Sigma)$ and $\bar{S}_q(\Sigma)$ hinges on the $\delta$-ISS-M$_q$ property. Given an input $u \in \mathcal{U}$, all solution processes of $\Sigma$ under the input $u$ forget the mismatch between their initial conditions and converge to each other with respect to the $q$th moment metric. Therefore, the longer the applied inputs are, the less relevant is the mismatch between initial conditions. Then, the fundamental idea of the introduced abstractions consists in taking the $N$ consequive applied inputs as the state of the symbolic model.

The control synthesis over $\bar{S}_q(\Sigma)$ (resp. $\bar{S}_q(\Sigma)$) is simple as the outputs are deterministic points, whereas for $S_q(\Sigma)$ it is perhaps less intuitive. Hence, we discuss it in more details later in Subsection 5.3.

**Example 2** An example of an abstraction $S_q(\Sigma)$ with $N = 3$ and $U_q = \{0, 1\}$ is depicted in Figure 1, where the initial states are shown as targets of sourceless arrows. Regardless of the size of the state set and of its dimension, $S_q(\Sigma)$ only has eight possible states, namely:

$$X_q = \{(0,0,0),(0,0,1),(0,1,0),(0,1,1),(1,0,0),(1,0,1),(1,1,0),(1,1,1)\}.$$

In order to obtain some of the main results of this work, we raise an assumption on the $\delta$-ISS-M$_q$ Lyapunov function $V$ which we will work with, as follows:

$$|V(x,y) - V(x,z)| \leq \sqrt{q}||y-z||,$$  \hspace{1cm} (7)
for any \( x, y, z \in \mathbb{R}^n \), and for some concave function \( \tilde{g} \in \mathcal{H}_c \). The function \( \tilde{g} \) in (7) can be easily computed as long as one is interested to work in a compact subset of \( \mathbb{R}^n \). Particularly, for all \( x, y, z \in D \), where \( D \subseteq \mathbb{R}^n \) is compact, one can apply the mean value theorem to the function \( y \to V(x, y) \) to get

\[
|V(x, y) - V(x, z)| \leq \tilde{g}(\|y - z\|), \quad \text{where} \quad \tilde{g}(r) = \left( \max_{x, y \in D \setminus \Delta} \left\| \frac{\partial V(x, y)}{\partial y} \right\| \right) r.
\]

In particular, for the \( \delta\text{-ISS-M}_1 \) Lyapunov function \( V(x, x') := \sqrt{(x - x')^T P(x - x')} \), for some positive definite matrix \( P \in \mathbb{R}^{n \times n} \) and for all \( x, x' \in \mathbb{R}^n \), one obtains \( \tilde{g}(r) = \sqrt{\lambda_{\text{max}}(P)} r \) due to the triangle inequality, which satisfies (7) globally on \( \mathbb{R}^n \). Remark that for deterministic control systems, the concavity assumption of \( \tilde{g} \) is not required.

Before providing the main results of the paper, we need the following technical results.

**Lemma 2** Consider a stochastic control system \( \Sigma \), admitting a \( \delta\text{-ISS-M}_q \) Lyapunov function \( V \), and consider its corresponding symbolic model \( \mathcal{S}_q(\Sigma) \). We have that

\[
\eta \leq \left( \alpha^{-1} \left( e^{-\kappa V r} \max_{u_q \in U_q} \mathbb{E} \left[ V \left( \tilde{x}_{\tau q_u}(\tau), \tilde{x}_s \right) \right] \right) \right)^{1/q}, \quad (8)
\]

where

\[
\eta := \max_{u_q \in U_q, q_u \in \mathcal{X}_q} \max_{x_q' \in \text{Post}_{u_q}(x_q)} \left\| \tilde{x}_{\tau q_u}(\tau) - H_q(x_q') \right\|. \quad (9)
\]

The proof of Lemma 2 is provided in the Appendix. The next lemma provides similar result as the one in Lemma 2, but without explicitly using any Lyapunov function.

**Lemma 3** Consider a \( \delta\text{-ISS-M}_q \) stochastic control system \( \Sigma \) and its corresponding symbolic model \( \mathcal{S}_q(\Sigma) \). We have:

\[
\eta \leq \left( \beta \left( \max_{u_q \in U_q} \left\| \tilde{x}_{\tau q_u}(\tau) - x_s \right\|^q N \tau \right) \right)^{1/q}, \quad (10)
\]

where \( \eta \) is given in (9) and \( \beta \) is the \( \mathcal{H}_c \mathcal{L} \) function appearing in (2).

The proof of Lemma 3 is provided in the Appendix. The next two lemmas provide similar results as Lemmas 2 and 3, but by using the symbolic model \( S_q(\Sigma) \) rather than \( \mathcal{S}_q(\Sigma) \).

**Lemma 4** Consider a stochastic control system \( \Sigma \), admitting a \( \delta\text{-ISS-M}_q \) Lyapunov function \( V \), and consider its corresponding symbolic model \( S_q(\Sigma) \). One has:

\[
\tilde{\eta} \leq \left( \alpha^{-1} \left( e^{-\kappa V r} \max_{u_q \in U_q} \mathbb{E} \left[ V \left( \tilde{x}_{\tau q_u}(\tau), \tilde{x}_s \right) \right] \right) \right)^{1/q}, \quad (11)
\]

where

\[
\tilde{\eta} := \max_{u_q \in U_q, q_u \in \mathcal{X}_q} \mathbb{E} \left( \left\| \tilde{z}_{H_q(x_{q_u})q_u}(\tau) - H_q(x_q') \right\| \right). \quad (12)
\]
Proof The proof is similar to the one of Lemma 2 and can be shown by using convexity of $\alpha$ and Jensen inequality [16].

**Lemma 5** Consider a $\delta$-ISS-$M_q$ stochastic control system $\Sigma$ and its corresponding symbolic model $S_q(\Sigma)$. We have:

$$\hat{\eta} \leq (\beta \left( \max_{u \in U_q} \mathbb{E} \left[ \left\| \xi_{u_0 u}(\tau) - x_\tau \right\| \right] . N \tau \right)^{1/q},$$

where $\hat{\eta}$ is given in (12) and $\beta$ is the $\mathcal{X} \mathcal{L}$ function appearing in (2).

**Proof** The proof is similar to the one of Lemma 3 and can be shown by using Jensen inequality [16].

**Remark 2** It can be readily verified that by choosing $N$ sufficiently large, $\eta$ and $\hat{\eta}$ can be made arbitrarily small. One can as well try to reduce the upper bound for $\eta$ (in (8) for example) by selecting the source state $x_\tau$ as follows:

$$x_\tau = \arg \min_{x \in \mathbb{R}^m; q^0 = \tau} \max_{x \in \mathbb{R}^m} V(\xi_{x u}(\tau), x).$$

(14)

We can now present the first main result of the paper, which relates the existence of a $\delta$-ISS-$M_q$ Lyapunov function to the construction of an approximately bisimilar symbolic model.

**Theorem 2** Consider a stochastic control system $\Sigma$ with $f(0_n, 0_m) = 0_n$ and $\sigma(0_n) = 0_{n \times p}$, admitting a $\delta$-ISS-$M_q$ Lyapunov function $V$, of the form of the one explained in Lemma 1, such that (7) holds for some concave $\hat{\gamma} \in \mathcal{K}_{\infty}$. Let $\eta$ be given by (9). For any $\epsilon \in \mathbb{R}^+$ and any tuple $q = (\tau, \mu, N, x_\tau)$ of parameters satisfying $\mu \leq \text{span}(U)$ and $N$ sufficiently large and

$$e^{-\kappa \tau} \alpha(\epsilon^q) + \frac{1}{\epsilon \kappa} \rho(\mu) + \hat{\gamma}\left(h_\mu(\left(\left(N + 1\right)\tau\right)) + \eta\right) \leq \alpha(\epsilon^q),$$

the relation (cf. Definition 5)

$$R = \left\{ (x_\tau, x_q) \in X_\tau \times X_q \mid \mathbb{E} \left[ V(x_\tau, \bar{H}_q(x_q)) \right] \leq \alpha(\epsilon^q) \right\}$$

is an $\epsilon$-approximate bisimulation relation between $S_q(\Sigma)$ and $S_\tau(\Sigma)$.

The proof can be found in the Appendix. By choosing $N$ sufficiently large and using the results in Lemmas 1 and 2, one can enforce $h_\mu(\left(\left(N + 1\right)\tau\right))$ and $\eta$ in (15) to be sufficiently small. Hence, for a given precision $\epsilon \in \mathbb{R}_+^+$, there always exists a sufficiently small value of $\mu$ and a large value of $N$, such that the condition in (15) is satisfied. A result similar as that in Theorem 2 can be recovered for a $\delta$-ISS deterministic control system $\Sigma$, as provided in the following corollary.

**Corollary 1** Consider a deterministic control system $\Sigma$ admitting a $\delta$-ISS Lyapunov function $V$ such that (7) holds for some $\hat{\gamma} \in \mathcal{K}_{\infty}$. Let $\eta$ be given by (9). For any $\epsilon \in \mathbb{R}^+$ and any tuple $q = (\tau, \mu, N, x_\tau)$ of parameters satisfying $\mu \leq \text{span}(U)$ and

$$e^{-\kappa \tau} \alpha(\epsilon) + \frac{1}{\epsilon \kappa} \rho(\mu) + \hat{\gamma}(\eta) \leq \alpha(\epsilon),$$

the relation

$$R = \left\{ (x_\tau, x_q) \in X_\tau \times X_q \mid V(x_\tau, \bar{H}_q(x_q)) \leq \alpha(\epsilon) \right\}$$

is an $\epsilon$-approximate bisimulation relation between $S_q(\Sigma)$ and $S_\tau(\Sigma)$.
The proof is similar to the one of Theorem 2. In order to mitigate the conservativeness that might rise from using Lyapunov functions, the next theorem provides a result that is similar to the one of Theorem 2, which is however not obtained by explicit use of $\delta$-ISS-M$_q$ Lyapunov functions, but by using functions $\beta$ and $\gamma$ as in (2).

**Theorem 3** Consider a $\delta$-ISS-M$_q$ stochastic control system $\Sigma$, satisfying the result of Lemma 1. Let $\eta$ be given by (9). For any $\varepsilon \in \mathbb{R}^+$, and any tuple $q = (\tau, \mu, N, x_s)$ of parameters satisfying $\mu \leq \text{span}(U)$ and

$$\left(\beta(\varepsilon^q, \tau) + \gamma(\mu)\right)^{\frac{1}{2}} + h_x((N + 1)\tau)^{\frac{1}{2}} + \eta \leq \varepsilon,$$

(17)

the relation

$$R = \left\{ (x_{\tau}, x_q) \in X_{\tau} \times X_q \mid \left\| x_{\tau} - H_x(x_q) \right\| \leq \left(\text{span}(U) \right)^{\frac{1}{2}} \right\}$$

is an $\varepsilon$-approximate bisimulation relation between $\mathcal{S}_q(\Sigma)$ and $\mathcal{S}_\tau(\Sigma)$.  

The proof can be found in the Appendix. By choosing $N$ sufficiently large and using the results in Lemmas 1 and 3, one can force $h_x((N + 1)\tau)$ and $\eta$ in (17) to be sufficiently small. Hence, for a given precision $\varepsilon$, there always exist a sufficiently large value of $\tau$ and $N$ and a small enough value of $\mu$ such that the condition in (17) is satisfied. However, unlike the result in Theorem 2, here for a given fixed sampling time $\tau$, one may not find any values of $N$ and $\mu$ satisfying (17) because the quantity $(\beta(\varepsilon^q, \tau))^{\frac{1}{2}}$ may be larger than $\varepsilon$. The symbolic model $\mathcal{S}_q(\Sigma)$, computed using the parameter $q$ provided in Theorem 3 (whenever existing), is likely to have fewer states than the one computed using the parameter $q$ provided in Theorem 2 – a similar fact has been experienced in the first example in [31]. A result similar to the one in Theorem 3 can be fully recovered for a $\delta$-ISS deterministic control system $\Sigma$, as provided in the following corollary.

**Corollary 2** Consider a $\delta$-ISS deterministic control system $\Sigma$. Let $\eta$ be given by (9). For any $\varepsilon \in \mathbb{R}^+$, and any tuple $q = (\tau, \mu, N, x_s)$ of parameters satisfying $\mu \leq \text{span}(U)$ and

$$\beta(\varepsilon, \tau) + \gamma(\mu) + \eta \leq \varepsilon,$$

(18)

the relation

$$R = \left\{ (x_{\tau}, x_q) \in X_{\tau} \times X_q \mid \left\| x_{\tau} - H_x(x_q) \right\| \leq \varepsilon \right\}$$

is an $\varepsilon$-approximate bisimulation relation between $\mathcal{S}_q(\Sigma)$ and $\mathcal{S}_\tau(\Sigma)$.  

The proof is similar to the one of Theorem 3. The next two theorems provide results that are similar to those of Theorems 2 and 3, but by using the symbolic model $\mathcal{S}_q(\Sigma)$ rather than $\mathcal{S}_q(\Sigma)$.

---

3 Here, $\beta$ and $\gamma$ are the $\mathcal{K}_L$ and $\mathcal{K}_\infty$ functions, respectively, appearing in (3).
Finally, in the conditions (15), (16), (17), (18), (19), and (20) set $\mathcal{U} = \emptyset$. For any $\varepsilon \in \mathbb{R}^+$, and any tuple $d = (\tau, \mu, N, x_0)$ of parameters satisfying $\mu \leq \text{span}(\mathcal{U})$ and

$$e^{-\kappa \mathcal{A}(\varepsilon^\theta)} + \frac{1}{e\kappa} \mathcal{D}(\mu) + \tilde{\gamma}(\tilde{\eta}) \leq \mathcal{A}(\varepsilon^\theta),$$

(19)

the relation

$$R = \{(x_t, x_0) \in \mathcal{X}_\tau \times \mathcal{X}_0 \mid \mathbb{E}[V(x_t, H_0(x_0))] \leq \mathcal{A}(\varepsilon^\theta)\}$$

is an $\varepsilon$-approximate bisimulation relation between $S_q(\Sigma)$ and $S_\tau(\Sigma)$. □

The proof is similar to the one of Theorem 2.

**Theorem 5** Consider a $\delta$-ISS-M$_\eta$ stochastic control system $\Sigma$. Let $\tilde{\eta}$ be given by (12). For any $\varepsilon \in \mathbb{R}^+$, and any tuple $d = (\tau, \mu, N, x_0)$ of parameters satisfying $\mu \leq \text{span}(\mathcal{U})$ and

$$(\beta(\varepsilon^\tau, \tau) + \gamma(\mu))^{\frac{1}{\tau}} + \tilde{\eta} \leq \varepsilon,$$

(20)

the relation

$$R = \{(x_t, x_0) \in \mathcal{X}_\tau \times \mathcal{X}_0 \mid \mathbb{E}[(\|x_t - H_0(x_0)||^\eta)]^{\frac{1}{\eta}} \leq \varepsilon\}$$

is an $\varepsilon$-approximate bisimulation relation between $S_q(\Sigma)$ and $S_\tau(\Sigma)$. □

The proof is similar to the one of Theorem 3.

**Remark 3** The symbolic model $S_q(\Sigma)$, proposed in Theorem 4 (resp. Theorem 5), has fewer (or at most equal number of) states than the symbolic model $\mathcal{S}_q(\Sigma)$, proposed in Theorem 2 (resp. Theorem 3) while having the same precision. However, the states in $S_q(\Sigma)$ have probabilistic output values, rather than deterministic ones, which is likely to make control synthesis procedures more involved (cf. Subsection 5.3). □

**Remark 4** Although we assume that the set $\mathcal{U}$ is infinite, Theorems 2, 3, 4, and 5 and Corollaries 1 and 2 still hold when the set $\mathcal{U}$ is finite, with the following modifications. First, the systems $\Sigma$ and $\mathcal{S}$ are required to satisfy properties (2) and (3), respectively, for $v = v'$. Second, take $U_q = U$ in the definitions of $\mathcal{S}_q(\Sigma)$ (resp. $S_q(\Sigma)$) and $\mathcal{S}_\tau(\Sigma)$. Finally, in the conditions (15), (16), (17), (18), (19), and (20) set $\mu = 0$. □

Finally, we establish the results on the existence of symbolic models $S_q(\Sigma)$ (resp. $S_q(\Sigma)$) and $\mathcal{S}_q(\Sigma)$ such that $S_q(\Sigma) \cong S_\tau(\Sigma)$ (resp. $S_q(\Sigma) \cong S_\tau(\Sigma)$) and $\mathcal{S}_q(\Sigma) \cong \mathcal{S}_\tau(\Sigma)$.

**Theorem 6** Consider the results in Theorem 2. If we select

$$X_{e_0} = \left\{ x \in \mathbb{R}^n \mid \|x - \mathcal{H}_q(x_{q_0})\| \leq \left(\mathcal{P}^{-1}(\mathcal{A}(\varepsilon^\theta))\right)^{\frac{1}{\tau}}, \exists x_{q_0} \in \mathcal{X}_{q_0} \right\},$$

then we have $S_q(\Sigma) \cong S_\tau(\Sigma)$. □
Theorem 8 Consider the results in Theorem 4. Let $\mathcal{A}$ denote the set of all $\mathbb{R}^n$-valued random variables, measurable over $\mathcal{F}_0$. If we select

$$X_{x_0} = \left\{ a \in \mathcal{A} \mid \mathbb{E} \left[ \| a - H_q(x_{q_0}) \| \right] \leq \sqrt[\frac{1}{2}]{{\alpha}^{-1} \left( \frac{\alpha}{\epsilon} \right)} \right\}, \exists x_{q_0} \in X_{q_0} \right\},$$

then we have $S_q(\Sigma) \preceq_{\epsilon, \gamma} S(\Sigma).$
The proof is similar to the one of Theorem 6.

**Theorem 9** Consider the results in Theorem 5. Let $\mathcal{A}$ denote the set of all $\mathbb{R}^n$-valued random variables, measurable over $\mathcal{F}_0$. If we select

$$X_{t0} = \left\{ a \in \mathcal{A} \mid (E[\|a - H_q(x_{q0})\|_q])^{1/q} \leq \varepsilon, \exists x_{q0} \in X_{q0} \right\},$$

then we have $S_q(\Sigma) \cong_{\varepsilon/\rho} S_r(\Sigma)$. \hfill $\Box$

The proof is similar to the one of Theorem 7.

### 5.3 Control synthesis over $S_q(\Sigma)$

Note that both $\overline{S}_q(\Sigma)$ and $S_q(\Sigma)$ are finite systems. The only difference is that the outputs of the former system are always deterministic points, whereas those of the latter can be non-degenerate random variables. Let us describe the control synthesis for these systems over quantitative specifications, and for example over the safety formula $\Box \varphi_A$, for $A \subseteq \mathbb{R}^n \subseteq Y$ (as already been used in Subsection 5.1). Clearly, since the original system $S_r(\Sigma)$ is stochastic in the sense that its outputs are non-degenerate random variables similarly to $S_q(\Sigma)$, it would be too conservative to require that it satisfies the formula exactly. Thus, we are rather interested in an input policy that makes $S_r(\Sigma)$ satisfy $\Box \varphi_A$, with some $\varepsilon > 0$: recall from Subsection 5.1 that the latter LTL formula can be satisfied by non-degenerate random variables, in contrast to $\Box \varphi_A$. Let us recap how to use abstractions for this task, and let us start with $S_q(\Sigma)$ belonging to a more familiar type of systems in which the outputs of states are deterministic.

We label a state $x_q$ of $S_q(\Sigma)$ with $A$ if $H_q(x_q) \in A$ and, say, with $B$ otherwise. As a result, we obtain a transition system with labels over the states and can synthesize a control strategy by solving a safety game [24] that makes an output run of $S_q(\Sigma)$ satisfy $\Box \varphi_A$. After that, we can exploit $\varepsilon$-approximate bisimilarity to guarantee that the refined input policy makes the corresponding output run of the original system $S_r(\Sigma)$ satisfy $\Box \varphi_A$.

The main subtlety in the case of $S_q(\Sigma)$ is how to label its states. We cannot do this as for $\overline{S}_q(\Sigma)$, since $H_q(x_q)$ may never be an element of $A$ for any $x_q \in X_q$: indeed, the latter is a set of deterministic points, whereas all the outputs of $S_q(\Sigma)$ can happen to be non-degenerate random variables. In order to cope with this issue, we propose to relax the original problem and at the same time to strengthen the quality of the abstraction. Namely, we can consider a relaxed problem $\Diamond \varphi_{A_\delta}$ over the abstraction $\overline{S}_q(\Sigma)$, for some $\delta \in [0, \varepsilon]$, where the latter is now required to be $(\varepsilon - \delta)$-approximate (rather than $\varepsilon$-approximate) bisimilar to the original system. Clearly $(A_{\delta})_{\varepsilon - \delta} \subseteq A_{\varepsilon}$, so that whenever the control policy for $\Diamond \varphi_{A_{\delta}}$ is synthesized over $\overline{S}_q(\Sigma)$, its refined version is guaranteed to enforce $\Diamond \varphi_{A_{\delta}}$ over the original system. Thanks to the fact that $A_{\delta}$ contains non-degenerate random variables, we eliminate the conservativeness presented before in the sense that it is likely that there are now points $x_q \in X_q$ in $S_q(\Sigma)$ such that $H_q(x_q) \in A_{\delta}$. The only remaining question is how to check whether $H_q(x_q) \in A_{\delta}$. To answer this question, we check that the distance

$$d(H_q(x_q), A) = \inf_{a \in A} (E[\|H_q(x_q)(N\tau) - a\|_q])^{1/q}$$

(21)
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is smaller than $\delta$, which involves both computing the expectation over the solution of the SDE, and optimizing the value of this expectation. Clearly, such a computation in general cannot be done analytically, and the evaluation of the expectation itself is a highly non-trivial task unless the SDE has a very special form.

We propose a Monte Carlo approach to compute an approximation of the quantity in (21) by means of empirical expectations. Using such an approach, we can estimate $d(H_q(x_0), A)$ only up to some precision, say $\theta$. If the estimated distance is less than $\delta - \theta$, we are safe to label $x_0$ with $A$, whereas all other states are labeled by $B$.

Furthermore, since this result is based on a Monte Carlo method, it holds true only with a certain confidence level $1 - \pi$ where $\pi \in [0, 1]$. The benefit of our approach is that it is not only valid asymptotically (as the number of samples grows to infinity), but we are also able to provide a number of simulations that is sufficient to estimate $d(H_q(x_0), A)$ with any given precision $\theta$ and with any given confidence $1 - \pi$. This can be considered as an extension of the well-known Hoeffding’s inequality [8] to the case when one has to deal with an optimization problem. Note that like in other cases when using Monte-Carlo for distance estimation, there are two levels of error. First, in $\varepsilon$ the $\varepsilon$ still refers to the precision of symbolic model in the $q$th moment metric, but now also a level of (lack of) confidence $\pi \in [0, 1]$ is added on top of it.

Regardless of the specification of interest, the main task over $S_q(\Sigma)$ is always to compute some distance as in (21) for any set that appears in the specification, so the method below applies not only to the safety formula $\square \varphi_A$, but also to more general formulae, which are left as object of the future research.

Suppose that $A$ as in (21) is a compact subset of $\mathbb{R}^n$, and let $A'$ be the smallest subset of $[\mathbb{R}^n]_r$ such that $A \subseteq \bigcup_{p \in A'} B^r_p \{p\}$. Let $M$ be the number of samples and let

$$d_M' := \min_{a \in A'} \left( \frac{1}{M} \sum_{i=1}^{M} \left\| \frac{\xi_i}{\Lambda_{x_0}} (N \tau) - a \right\|^q \right)^{\frac{1}{q}},$$

where the superscript $i$ denotes the index of samples. Now, we have the following theorem.

**Theorem 10** For any stochastic control system $\Sigma$ one has $|d(H_q(x_0), A) - d_M'| \leq \theta$ with confidence of at least $1 - \pi$, given that $r < 20\theta$ and that

$$M \geq \frac{|A'|b(a^*, 2q)}{\pi(\theta - r/2)^{2q}},$$

where $b(a, p) := (1 + |x_0 - a|^p)e^{\theta(p+1)\max\{L_x, Z\} N^2}$ and $a^* \in \arg\max_{a' \in A'} \| x_0 - a' \|$.

The proof can be found in the Appendix. Let us make some comments on Theorem 10. First of all, no matter how many distances one has to evaluate, one can always use the same samples $\xi_i$ and there is no need to generate new samples. Second, to the best of our knowledge, logarithmic bounds on $M$ (as per [11]) are not available in this general case due to the fact that we deal with an unbounded state space.
5.4 Relationship with existing results in the literature

Note that given any precision \( \varepsilon \in \mathbb{R}^+ \) and sampling time \( \tau \), one can always use the results in Theorem 6 to construct a symbolic model \( \mathcal{S}_q(\Sigma) \) that is \( \varepsilon \)-approximate bisimilar to \( \mathcal{S}_\tau(\Sigma) \) without any state set discretization, but only input set discretization. The results in Theorem 5.1 in [31] also provide symbolic models that are \( \varepsilon \)-approximate bisimilar to \( \mathcal{S}_\tau(\Sigma) \). However, the results in [31] require both state and input set discretization and cannot be applied for any sampling time \( \tau \) as long as:

\[
\text{convenient to use the proposed symbolic model here rather than the one proposed in [31]}
\]

where \( K \in \mathbb{R}_+ \) is the state space quantization parameter, called here \( \nu \), is the same as long as both use the same input set quantization parameter \( \mu \). The reason their precisions are approximately (rather than exactly) the same is because we use \( h_{\nu}(\sigma, (N + 1)\tau) \) in conditions (15) and (17) in this paper rather than \( h(\tau) = \sup_{x \in \Sigma} h_1(\tau) \) that is being used in conditions 5.4 and 5.14 in [31] for a compact set \( \mathcal{D} \subset \mathbb{R}^n \). By assuming that \( h_{\nu}(\sigma, (N + 1)\tau) \) and \( h(\tau) \) are much smaller than \( \eta \) and \( \nu \), respectively, or \( h_{\nu}(\sigma, (N + 1)\tau) \approx h(\tau) \), one should expect to get the same precisions for the symbolic models provided here and those provided in [31].

The number of states of the proposed symbolic model in this paper is \( \lvert \mathcal{U} \rvert^N \). Assume that we are interested in the dynamics of \( \Sigma \) on a compact set \( \mathcal{D} \subset \mathbb{R}^n \). Since the set of states of the proposed symbolic model in [31] is \( \mathcal{D}_\nu \), its size is \( \lvert \mathcal{D}_\nu \rvert = \frac{K}{\nu} \), where \( K \) is a positive constant proportional to the volume of \( \mathcal{D} \). Hence, it is more convenient to use the proposed symbolic model here rather than the one proposed in [31] as long as:

\[
\| \mathcal{U} \|^{N} \leq \frac{K}{(\alpha^{-1}(e^{-\kappa n\tau} \eta_0))^n}.
\]

Without loss of generality, one can assume that \( \alpha(r) = r \) for any \( r \in \mathbb{R}_+^+ \). Hence, for sufficiently large value of \( N \), the size of the proposed symbolic model here is smaller than the one proposed in [31] as long as:

\[
\| \mathcal{U} \|^{N} \leq \frac{K}{(\alpha^{-1}(e^{-\kappa n\tau} \eta_0))^n} 
\]

One can readily verify that if the state-space dimension of \( \Sigma \) is very large, i.e. \( n \gg 1 \), or \( \Sigma \) is “strongly” \( \delta \)-ISS-Md, i.e. \( K \gg 1 \), inequality (22) is most likely satisfied.

Finally, remark that the framework proposed in this paper lets us to construct less conservative symbolic models with probabilistic output values while the proposed one in [31] only provides symbolic models with deterministic output values.
6 Example

We show the effectiveness of the results presented in this work by constructing a bisimilar symbolic model for the model of a road network. The road is divided in 5 cells of 250 meters with 2 entries and 2 ways out, as depicted schematically in Figure 2. The model is borrowed from [5], explained in details in [27], however it is now affected by noise and described in continuous time.

The two entries are controlled by traffic lights, denoted by $u_1$ and $u_2$, that enable (green light) or not (red light) the vehicles to pass. In this model the length of a cell is in kilometres (0.25 km), and the flow speed of the vehicles is 70 kilometers per hour (km/h). Moreover, during the sampling time interval $\tau$, it is assumed that 6 vehicles pass the entry controlled by the light $u_1$, 8 vehicles pass the entry controlled by the light $u_2$, and one quarter of vehicles that leave cell 1 goes out on the first exit. We assume that both lights cannot be red at the same time. The model of $\Sigma$ is described by:

$$d\xi = (A\xi + Bu)dt + \xi dW_t,$$

where

$$A = 10^4 \times \begin{pmatrix} -0.0541 & 0 & 0 & 0 & 0 \\ 0.3224 & -0.1370 & 0 & 0 & 0 \\ -0.7636 & 0.3224 & -0.0541 & 0 & 0 \\ 2.1122 & -0.7636 & 0.1260 & -0.0541 & 0 \\ -6.2132 & 2.1122 & -0.2205 & 0.1260 & -0.0541 \end{pmatrix},$$

$$B = 10^4 \times \begin{pmatrix} 0.0696 & 0 \\ -0.2743 & 0 \\ 0.7075 & 0.0696 \\ -2.0081 & -0.0924 \\ 5.9802 & 0.1911 \end{pmatrix},$$

$U = \{u_0, u_1, u_2\} = \{[6 \ 8]^T, [6 \ 0]^T, [0 \ 8]^T\}$, and $\xi_i$ is the number of vehicles in cell $i$ of the road. We point out that $\mathcal{U}_\tau$ contains curves taking values in $U$. Since $U$ is finite, as explained in Remark 4, $\mu = 0$ is to be used in (15), (17), (19), and (20). Using LMI
Remark 5

By considering the deterministic control system

\[
P = 10^4 \times \begin{bmatrix}
76763.4393 & -2101.1583 & 3790.9182 & -155.6576 & -125.9871 \\
-2101.1583 & 10676.9437 & 1237.3552 & -86.6855 & 100.5718 \\
3790.9182 & 1237.3552 & 1823.02431 & 171.1549 & -71.1162 \\
-155.6576 & 86.6855 & 171.1549 & 229.2134 & -5.5649 \\
-125.9871 & 100.5718 & -71.1162 & -5.5649 & 33.3977
\end{bmatrix},
\]

satisfies conditions (i)-(iii) in Definition 3 with \( q = 2 \), \( \kappa = 300 \), \( \alpha(r) = \frac{1}{2} \lambda_{\text{min}}(P)r \), \( \beta(r) = \frac{1}{2} \lambda_{\text{mac}}(P)r \), \( \rho(r) = \frac{5\theta_0^t \|P\|}{2k} r^2 \), \( \forall r \in \mathbb{R}_0^+ \). Hence, \( \Sigma \) is \( \delta \)-ISS-M_2, equipped with the \( \delta \)-ISS-M_2 Lyapunov function \( V \). Using the results of Theorem 1, provided in [31], one gets that functions \( \beta(r,s) = \alpha^{-1}(\beta(r)e^{-\kappa r}) \) and \( \gamma(r) = \beta^{-1}(\frac{1}{\varepsilon \kappa} \rho(r)) \) satisfy property (2) for \( \Sigma \). We choose the source state as the one picked in [5], i.e. \( x_s = [3.8570 \ 3.3750 \ 3.3750 \ 8.5177 \ 8.5177]^T \).

For a given precision \( \varepsilon = 0.5 \) and fixed sampling time \( \tau = 0.00277 \) 
(10 sec), the parameter \( N \) for \( \bar{S}_q(\Sigma) \), based on inequality (15) in Theorem 2, is obtained as 14. Therefore, the resulting cardinality of the set of states for \( \bar{S}_q(\Sigma) \) is \( |U|^{14} = 314 = 4782969 \). Using the aforementioned parameters, one gets \( \eta \leq 6.0776 \times 10^{-6} \), where \( \eta \) is given in (9). Remark that the results in Theorems 3 and 5 cannot be applied here because \( (\beta(e^q, \tau))^\frac{1}{2} \geq \varepsilon \). Using criterion (22), one has \( |U|e^{-\frac{1}{2\varepsilon}} = 0.37 \), implying that the proposed approach in this paper is more efficient in terms of the size of the abstraction than the one proposed in [31]. We elaborate more on this at the end of the section.

Remark 5

By considering the deterministic control system \( \Sigma \) and using the results in Corollary 1 and the same parameters \( q \) as the ones in \( \bar{S}_q(\Sigma) \), one obtains \( \varepsilon = 0.01 \) in (16). As expected, \( \bar{S}_q(\Sigma) \) (i.e. symbolic model for the deterministic control system \( \Sigma \)) provides much smaller precision than \( \bar{S}_q(\Sigma) \) (i.e. symbolic model for the stochastic control system \( \Sigma \)) while having the same size as \( \bar{S}_q(\Sigma) \).

Now the objective, as inspired by the one suggested in [5], is to design a schedule for the coordination of traffic lights enforcing \( \Sigma \) to satisfy a safety and a fairness property. The safety part is to keep the density of traffic lower than 16 vehicles per cell which can be encoded via the LTL specification \( \square \phi_W \), where \( W = [0 \ 16]^3 \) and \( \phi_W \) is a label characterizing the set \( W \). The fairness part requires to alternate the accesses between the two traffic lights and to allow only 3 identical consecutive modes of light ensuring fairness between two traffic lights. We implemented the proposed techniques in this paper on top of the recently developed synthesis toolbox SCOTS [22] on an iMac with CPU 3.5GHz Intel Core i7. The CPU time used for computing the abstraction and synthesizing the controller have amounted to 68 and 5 seconds, respectively. Starting from the initial condition \( x_0 = [1.417 \ 4.993 \ 10.962 \ 9.791 \ 14.734]^T \), we obtain a periodic schedule \( u = (u_0u_0u_0u_2u_1u_0u_0u_2u_1u_0u_0u_2u_1u_2)^6 \) keeping \( u_0 \) as much as possible in each period in order to maximize number of vehicles accessing the road. One can readily verify from the computed schedule \( u \) that the fairness property is satisfied.
Figure 3 displays a few realizations of the closed-loop solution process $\xi_{x_0,\nu}$. In Figure 3 bottom right, we show the average value (over 100000 experiments) of the distance (in the 2nd moment metric) in time of the solution process $\xi_{x_0,\nu}$ to the set $W$, namely $\|\xi_{x_0,\nu}(t)\|_{W}$, where the point-to-set distance is defined as $\|x\|_{W} = \inf_{w \in W} \|x - w\|$. Observe that the empirical average distance is as expected lower than the precision $\epsilon = 0.5$ due to the conservative nature of $\delta$-ISS-M function $V$.

![Fig. 3](image)

Fig. 3 A few realizations of the closed-loop solution process $\xi_{x_0,\nu}$ (top panel and the first two figures from the left in the bottom panel) and the average values (over 100000 experiments) of the distance of the solution process $\xi_{x_0,\nu}$ to the set $W$ (bottom right panel).

To compute exactly the size of the symbolic model, proposed in Theorem 5.1 in [31], we consider the dynamics of $\Sigma$ over the subset $D = [0 \ 16]^5$ of $\mathbb{R}^5$. Remark that Theorem 5.3 in [31] cannot be applied here because $(\beta(e^{q}, \tau))^{1/2} > \epsilon$. Using the same precision $\epsilon = 0.5$ and sampling time $\tau = 0.00277$ as the ones here, and the inequalities (5.3) and (5.4) in [31], we obtain the state space quantization parameter as $\nu \leq 6.0776 \times 10^{-6}$. Therefore, if one uses $\nu = 6.0776 \times 10^{-6}$, the cardinality of the state set of the symbolic model, provided by the results in Theorem 5.1 in [31], is equal to $(\frac{16}{\nu^5}) = 1.2645 \times 10^{32}$ which is practically not tractable and much higher than the one proposed here, amounting instead to 4782969 states.

7 Conclusions

In this paper we have proposed a symbolic abstraction technique for incrementally stable stochastic control systems (and corresponding deterministic model), which features only the discretization of the input set. The proposed approach is potentially more scalable than the one proposed in [31] for higher dimensional stochastic control systems.
Future work will concentrate on efficient implementations of the symbolic models proposed in this work on top of the recently developed synthesis toolbox SCOTS [22], as well as on more efficient controller synthesis techniques by using binary search trees data structure. The current implementation is not computationally very efficient for the closed-loop implementation of the synthesized symbolic controllers for deterministic control systems because one requires to loop over the symbolic states to find the ones with the outputs $\varepsilon$-close to the current measured state of the concrete system in each feedback iteration.

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Appendix

Proof of Lemma 2: Let \( x_q \in X_q \), where \( X_q = (u_1, u_2, \ldots, u_N) \), and \( u_q \in U_q \). Using the definition of \( \bar{T}_q(\Sigma) \), one obtains \( \tilde{x}_q = (u_2, \ldots, u_N, u_q) \in \text{Post}_{\tilde{u}_q}(x_q) \). Since \( V \) is a \( \delta \)-ISS-M \( \bar{T}_q \) Lyapunov function for \( \Sigma \), we have:

\[
\alpha \left( \left\| \bar{T}_q(x_q)(\tau) - \bar{T}_q(x_q') \right\| \right)^q 
\leq V(\bar{T}_q(x_q)(\tau), \bar{T}_q(x_q'))
\]

\[
= V(\bar{T}_q(x_q(N\tau), \bar{T}_q(x_q'))(N\tau), \bar{T}_q(x_q, u_2, \ldots, u_N)(N\tau))
\]

\[
\leq e^{-\kappa N\tau} V(\bar{T}_q(x_q(\tau), x_q)).
\] (24)

We refer the interested readers to the proof of Theorem 1 in [31] to see how we derived the inequality (24). Hence, one gets

\[
\left\| \bar{T}_q(x_q)(\tau) - \bar{T}_q(x_q') \right\| \leq (\alpha^{-1}(e^{-\kappa N\tau} V(\bar{T}_q(x_q(\tau), x_q))))^{1/q},
\] (25)

because of \( \alpha \in \mathcal{K}_\infty \). Since the inequality (25) holds for all \( x_q \in X_q \) and \( u_q \in U_q \), and \( \alpha \in \mathcal{K}_\infty \), inequality (8) holds.
Proof of Lemma 3: Let \(x_q \in X_q\), where \(x_q = (u_1, u_2, \ldots, u_N)\), and \(u_q \in U_q\). Using the definition of \(S_q(\Sigma)\), one obtains \(x_q' = (u_2, \ldots, u_N, u_q) \in \text{Post}_{u_q}(x_q)\). Since \(\Sigma\) is \(\delta\)-ISS-M and using inequality (2), we have:

\[
\|\hat{Q}_{u_q}(x_q) - \mathcal{H}_q(\tau)\|^q = \|\hat{Q}_{u_q}(N\tau) - \mathcal{H}_q(\tau)\|^q
\]

\[
= \|\hat{Q}_{u_q}(\tau(u_2, \ldots, u_N, u_q))(N\tau) - \hat{Q}_{u_q}(\tau)\|^q \leq \beta(\|\hat{Q}_{u_q}(\tau) - x_q\|^q, N\tau).
\]

Hence, one gets

\[
\|\hat{Q}_{u_q}(x_q) - \mathcal{H}_q(\tau')\| \leq (\beta(\|\hat{Q}_{u_q}(\tau) - x_q\|^q, N\tau))^{1/q}. \tag{26}
\]

Since the inequality (26) holds for all \(x_q \in X_q\) and all \(u_q \in U_q\), and \(\beta\) is a \(\mathcal{K}_\alpha\) function with respect to its first argument when the second one is fixed, inequality (10) holds.

Proof of Theorem 2: We start by proving that \(R\) is an \(\epsilon\)-approximate simulation relation from \(S_\Sigma(\Sigma)\) to \(S_q(\Sigma)\). Consider any \((\tau, x_q) \in R\). Condition (i) in Definition 5 is satisfied because

\[
(\mathbb{E}[\|x_\tau - \mathcal{H}_q(x_q)\|^q])^{1/q} \leq (\alpha^{-1}(\mathbb{E}[V(\tau, x_\tau)])^{1/q} \leq \epsilon. \tag{27}
\]

We used the convexity assumption of \(\alpha\) and the Jensen inequality [16] to show the inequalities in (27). Let us now show that condition (ii) in Definition 5 holds. Consider any \(u_\tau \in U_\tau\). Choose an input \(u_q \in U_q\) satisfying

\[
\|u_\tau - u_q\| = \|u_\tau(0) - u_q(0)\| \leq \mu. \tag{28}
\]

Note that the existence of such \(u_q\) is guaranteed by \(U\) being a finite union of boxes and by the inequality \(\mu \leq \text{span}(U)\) which guarantees that \(U \subseteq \bigcup_{\mu \in \mathbb{R}_\mu} \delta(\mu)\). Consider the transition \(x_\tau \xrightarrow{\tau} x_\tau' = x_\tau + u_\tau\) \(\mathcal{P}\)-a.s. in \(S_\Sigma(\Sigma)\). Since \(V\) is a \(\delta\)-ISS-M Lyapunov function for \(\Sigma\) and using inequality (28), we have (cf. equation (3.3) in [31])

\[
\mathbb{E}[V(x_\tau', \hat{Q}_{u_q}(x_q))] \leq \mathbb{E}[V(x_\tau, \mathcal{H}_q(x_q))]e^{-\alpha\tau} + \frac{1}{eK} \rho(\|u_\tau - u_q\|) \leq \alpha(e^\mu) e^{-\alpha\tau} + \frac{1}{eK} \rho(\mu). \tag{29}
\]

Observe that existence of \(u_q\), by the definition of \(S_q(\Sigma)\), implies the existence of \(x_q \xrightarrow{\tau} x_q'\) in \(S_q(\Sigma)\). Using Lemma 1, the concavity of \(\hat{Q}\), the Jensen inequality [16],
equation (9), the inequalities (7), (15), (29), and triangle inequality, we obtain
\[
E[V(x'_t, \xi_{\mathcal{P}}(q'(x'_q))) = E[V(x'_t, \xi_{\mathcal{P}}(q(x'_q))) + V(x'_t, \mathcal{H}_q(q(x'_q))) - V(x'_t, \xi_{\mathcal{P}}(q(x'_q)))] \\
= E[V(x'_t, \xi_{\mathcal{P}}(q(x'_q))) + E[V(x'_t, \mathcal{H}_q(q(x'_q))) - V(x'_t, \xi_{\mathcal{P}}(q(x'_q)))] \\
\leq \alpha(e^\delta) e^{-\gamma t} + \frac{1}{e^\kappa} \rho(\mu) + E[\tilde{\gamma}(\xi_{\mathcal{P}}(q(x'_q)))] \\
\leq \alpha(e^\delta) e^{-\gamma t} + \frac{1}{e^\kappa} \rho(\mu) \\
\quad + \tilde{\gamma}(\xi_{\mathcal{P}}(q(x'_q))) + ||\xi_{\mathcal{P}}(q(x'_q)) - \mathcal{H}_q(x'_q)||] \\
\leq \alpha(e^\delta) e^{-\gamma t} + \frac{1}{e^\kappa} \rho(\mu) + \tilde{\gamma}(h_q((N+1)t)\frac{1}{t} + \eta) \leq \alpha(e^\delta).
\]
Therefore, we conclude that \((x'_t, x'_q) \in R\) and that condition (ii) in Definition 5 holds.

Now we prove that \(R^{-1}\) is an \(\varepsilon\)-approximate simulation relation from \(S_t(\Sigma)\) to \(S_t(\Sigma)\). Consider any \((x_t, x_q) \in R\) (or equivalently \((x_q, x_t) \in R^{-1}\)). As showed in the first part of the proof, condition (i) in Definition 5 is satisfied. Let us now show that condition (ii) in Definition 5 holds. Consider any \(u_q \in U_q\). Choose the input \(u_t = u_q\) and consider \(x'_t = \xi_{\mathcal{P}}(\xi_{\mathcal{P}}(q(x'_q)))\) \(P\)-a.s. in \(S_t(\Sigma)\). Since \(V\) is a ISS-M \(\delta\) Lyapunov function for \(\Sigma\), one obtains (cf. equation 3.3 in [31]):
\[
E[V(x'_t, \xi_{\mathcal{P}}(q(x'_q)))] \leq e^{-\gamma t} E[V(x_t, \mathcal{H}_q(q(x'_q)))] \leq e^{-\gamma t} \alpha(e^\delta).
\]
Using Lemma 1, the definition of \(S_t(\Sigma)\), the concavity of \(\tilde{\gamma}\), the Jensen inequality [16], equation (9), the inequalities (7), (15), (30), and triangle inequality, we obtain
\[
E[V(x'_t, \mathcal{H}_q(q(x'_q)))] = E[V(x'_t, \xi_{\mathcal{P}}(q(x'_q)) + V(x'_t, \mathcal{H}_q(q(x'_q)) - V(x'_t, \xi_{\mathcal{P}}(q(x'_q)))] \\
= E[V(x'_t, \xi_{\mathcal{P}}(q(x'_q)))] + E[V(x'_t, \mathcal{H}_q(q(x'_q))) - V(x'_t, \xi_{\mathcal{P}}(q(x'_q)))] \\
\leq e^{-\gamma t} \alpha(e^\delta) + E[\tilde{\gamma}(\xi_{\mathcal{P}}(q(x'_q)))] \\
\leq e^{-\gamma t} \alpha(e^\delta) + \tilde{\gamma}(\xi_{\mathcal{P}}(q(x'_q))) + ||\xi_{\mathcal{P}}(q(x'_q)) - \mathcal{H}_q(x'_q)||] \\
\leq e^{-\gamma t} \alpha(e^\delta) + \tilde{\gamma}(\xi_{\mathcal{P}}(q(x'_q))) + ||\xi_{\mathcal{P}}(q(x'_q)) - \mathcal{H}_q(x'_q)||] \\
\leq e^{-\gamma t} \alpha(e^\delta) + \tilde{\gamma}(h_q((N+1)t)\frac{1}{t} + \eta) \leq \alpha(e^\delta).
\]
Therefore, we conclude that \((x'_t, x'_q) \in R\) (or equivalently \((x'_q, x'_t) \in R^{-1}\)) and condition (ii) in Definition 5 holds.

Proof of Theorem 3: We start by proving that \((x'_t, x'_q) \in R\) and that condition (ii) in Definition 5 holds. Consider any \(u_t \in U_t\). Choose an input \(u_q \in U_q\) satisfying
\[
||u_t - u_q||_\infty = ||u_t(0) - u_q(0)|| \leq \mu.
\]

Note that the existence of such $u_q$ is guaranteed by $U$ being a finite union of boxes and by the inequality $\mu \leq \text{span}(U)$ which guarantees that $U \subseteq \bigcup_{\nu \in [0, \mu]} \mathcal{B}_\mu(p)$. Consider the transition $x_\tau \xrightarrow{u_\delta} x'_\tau = \xi_{x_\tau} u_{\tau} (\tau)$ $\mathbb{P}$-a.s. in $S_\tau(\Sigma)$. It follows from the $\delta$-ISS-M$_q$ assumption on $\Sigma$ and (31) that:

$$
\mathbb{E}[\|x'_\tau - \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau)\|^q] \leq \beta(\mathbb{E}[\|x_\tau - \mathcal{P}_q(x_\tau)\|^q], \tau) + \gamma(\|u_\tau - u_q\|_\infty) \leq \beta(\epsilon^q, \tau) + \gamma(\mu).
$$

(32)

Existence of $u_q$, by the definition of $\mathcal{S}_q(\Sigma)$, implies the existence of $x_\tau \xrightarrow{u_q} x'_\tau$ in $\mathcal{S}_q(\Sigma)$, Using equation (9), the inequalities (5), (17), (32), and triangle inequality, we obtain

$$
\begin{align*}
(\mathbb{E}[\|x'_\tau - \mathcal{P}_q(x'_\tau)\|^q])^{\frac{1}{q}} &= (\mathbb{E}[\|x'_\tau - \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) + \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) \mathcal{P}_q(x_\tau)\|^q])^{\frac{1}{q}} \\
&- \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) + \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) - \mathcal{P}_q(x'_\tau)\|_{\|q\|}\frac{1}{q} \\
&\leq (\mathbb{E}[\|x'_\tau - \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau)\|^q])^{\frac{1}{q}} + (\mathbb{E}[\|\xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) - \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau)\|^q])^{\frac{1}{q}} \\
&+ (\mathbb{E}[\|\xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) - \mathcal{P}_q(x'_\tau)\|_{\|q\|}\frac{1}{q})^{\frac{1}{q}} \\
&\leq (\beta(\epsilon^q, \tau) + \gamma(\mu))^{\frac{1}{q}} + (h_{\xi}((N + 1)\tau))^\frac{1}{q} + \eta \leq \epsilon. 
\end{align*}
$$

Therefore, we conclude that $(x'_\tau, x'_\tau) \in R$ and that condition (ii) in Definition 5 holds.

Now, we prove that $R^{-1}$ is an $\epsilon$-approximate simulation relation from $\mathcal{S}_q(\Sigma)$ to $S_\tau(\Sigma)$. Consider any $(x_\tau, x_{\tau}) \in R$ (or equivalently $(x_{\tau}, x_\tau) \in R^{-1}$). Condition (i) in Definition 5 is satisfied by the definition of $R$. Let us now show that condition (ii) in Definition 5 holds. Consider any $u_q \in U_q$. Choose the input $u_\tau = u_q$ and consider $x'_\tau = \xi_{x_\tau} u_{\tau}(\tau)$ $\mathbb{P}$-a.s. in $S_\tau(\Sigma)$. Since $\Sigma$ is $\delta$-ISS-M$_q$, one obtains:

$$
\mathbb{E}[\|x'_\tau - \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau)\|^q] \leq \beta(\mathbb{E}[\|x_\tau - \mathcal{P}_q(x_\tau)\|^q], \tau) \leq \beta(\epsilon^q, \tau).
$$

(33)

Using definition of $\mathcal{S}_q(\Sigma)$, equation (9), the inequalities (5), (17), (33), and the triangle inequality, we obtain

$$
\begin{align*}
(\mathbb{E}[\|x'_\tau - \mathcal{P}_q(x'_\tau)\|^q])^{\frac{1}{q}} &= (\mathbb{E}[\|x'_\tau - \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) + \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) \mathcal{P}_q(x_\tau)\|^q])^{\frac{1}{q}} \\
&- \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) + \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) - \mathcal{P}_q(x'_\tau)\|_{\|q\|}\frac{1}{q} \\
&\leq (\mathbb{E}[\|x'_\tau - \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau)\|^q])^{\frac{1}{q}} + (\mathbb{E}[\|\xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) - \xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau)\|^q])^{\frac{1}{q}} \\
&+ (\mathbb{E}[\|\xi_{\mathcal{P}_q(x_\tau) u_{\tau}}(\tau) - \mathcal{P}_q(x'_\tau)\|_{\|q\|}\frac{1}{q})^{\frac{1}{q}} \\
&\leq (\beta(\epsilon^q, \tau) + (h_{\xi}((N + 1)\tau))^\frac{1}{q} + \eta \leq \epsilon. 
\end{align*}
$$

Therefore, we conclude that $(x'_\tau, x'_\tau) \in R$ (or equivalently $(x'_\tau, x'_\tau) \in R^{-1}$) and condition (ii) in Definition 5 holds. □
Proof of Theorem 10: Denote \( \hat{\theta} := \theta - r/2 > 0 \), and \( d_M(a) := \left( \frac{1}{M} \sum_{i=1}^{M} \| \xi_{t,a_q}^i - a \|^q \right)^{\frac{1}{q}} \) for all \( a \in \mathbb{R}^a \). It follows from [11, Theorem 4.5.4] that for all \( p \geq 1 \) and \( a \in \mathbb{R}^a \)

\[
\mathbb{E} \left[ \| \xi_{t,a_q} (\mathcal{N} \tau) - a \|^p \right] \leq b(a, p).
\]

Since we do not assume that the set of continuous states is bounded, the distance can be any positive real number, and the usual method of applying Hoeffding’s inequality does not work in this case. Instead we use Chernoff-type inequality (e.g. see above formula (1) in [4]), which implies that for any \( a' \in A' \) :

\[
P \left( \left| \| d(H_q(x_q), a') \| - d_M(a') \right| \geq \hat{\theta} \right) \leq \frac{b(a', 2q)}{M \hat{\theta}^2}.
\]

Furthermore, since \( x \mapsto x^q \) is Hölder continuous with power \( q \),

\[
P \left( \left| \| d(H_q(x_q), a') - d_M(a') \| \geq \hat{\theta} \right) \leq \frac{b(a', 2q)}{M \hat{\theta}^2 q}.
\]

Thus, for the union of such events over \( a' \in A' \), we have

\[
P \left( \exists a' \in A' \text{ s.t. } \| d(H_q(x_q), a') - d_M(a') \| \geq \hat{\theta} \right) \leq \frac{|A'| b(a'^{r}, 2q)}{M \hat{\theta}^{2q}}; \quad (34)
\]

due to the fact that the probability of a union is dominated by the sum of probabilities.

Let \( \lfloor \cdot \rfloor : A \rightarrow A' \) be any surjective map such that \( ||a - [a]|| \leq r/2 \) for all \( a \in A \), i.e. \( \lfloor \cdot \rfloor \) chooses an \( r/2 \)-close point in the grid \( A' \). Using this map, we can extrapolate the inequality (34) to the whole set \( A \) since

\[
\| d(H_q(x_q), a) - d_M([a]) \| \leq \| d(H_q(x_q), a) - d(H_q(x_q), [a]) \| + \| d(H_q(x_q), [a]) - d_M([a]) \| \\
\leq r/2 + \| d(H_q(x_q), [a]) - d_M([a]) \|,
\]

where we used the fact that \( \| d(H_q(x_q), a) - d(H_q(x_q), [a]) \| \leq ||a - [a]|| \) by the triangle inequality. As a result, the following inequality holds:

\[
P \left( \exists a \in A \text{ s.t. } \| d(H_q(x_q), a) - d_M([a]) \| \leq \theta \right) \leq p \left( \exists a' \in A' \text{ s.t. } \| d(H_q(x_q), a') - d_M(a') \| \geq \hat{\theta} \right). \quad (35)
\]

On the other hand, since for any two functions \( f, g : A \rightarrow \mathbb{R} \) it holds that

\[
\left| \inf_{a \in A} f(a) - \inf_{a \in A} g(a) \right| \leq \sup_{a \in A} |f(a) - g(a)|,
\]

we obtain that

\[
P \left( \| d(H_q(x_q), A) - d_M'^q \| \geq \theta \right) \leq P \left( \exists a \in A \text{ s.t. } \| d(H_q(x_q), a) - d_M([a]) \| \leq \theta \right).
\]

Combining the latter inequality with (34) and (35) yields:

\[
P \left( \| d(H_q(x_q), A) - d_M'^q \| \geq \theta \right) \leq \frac{|A'| b(a'^{r}, 2q)}{M \hat{\theta}^{2q}},
\]

and in case \( M \) satisfies the assumption of the theorem, the right-hand side is bounded above by \( \pi \) as desired. \( \square \)
Dear editor-in-chief, guest editor, and reviewers,

We would like to thank you for taking the time to handle and review our paper. We also appreciate the several suggestions and comments for improvements, which we have incorporated in the revised manuscript. Major changes are highlighted in blue color in the revised version. We present below a detailed discussion of all the questions and concerns raised by the guest editor and reviewers.

With kindest regards,

Majid Zamani, Ilya Tkachev, and Alessandro Abate
Individual responses

Response to the Guest Editor

There is some concern about the overlap between the current paper and references [26] and [28]. Please address the reviewer comments.

Response: We thank the guest editor for highlighting this concern raised by one of the reviewer: we address this concern in detail in a later subsection referring to the specific reviewer.

In addition, you might want to include some comparison with a few other works that use input/output discretization instead of state discretization such as:
* D. Tarraf, “An input-output construction of finite state ρ/μ approximations for control design”, TAC 2014.
* Schmuck et al., “Comparing Asynchronous l-Complete Approximations and Quotient Based Abstractions” arXiv 2015.

Response: We thank the guest editor for providing those references: we have provided a comparison with those articles in the introduction of the revised manuscript.

Response to Reviewer 1

Page 9: Here you start discussing LTL, but what is the LTL semantics for random variables? I do not think that this is standard for most readers of the journal.

Response: Since we describe the original stochastic control systems and its abstractions using non-blocking (in)finite transition systems with the trace-based semantics (cf. definition of output run after (6)) [3], the semantics of LTL is also trace-based and their interpretations over transition systems follow, correspondingly [3, Subsection 5.1.2]. We have prefered not to elaborate the syntax and semantic of LTL formulae formally in the paper, for the sake of brevity and because they are widely known and used by now. However, in the revised version we explicitly refer to [3], where LTL syntax and semantic are formally introduced.

Page 9, Line 19: Since you are not giving the semantics of LTL, it would be better to provide a simple explanation what □ϕA means to the reader.

Response: We thank the reviewer for the suggestion. We have explicitly explained what □ϕA means in the revised manuscript.

Page 9, Line 23: Why does this imply that y has a Dirac distribution? Doesn’t it just mean that the probability measure of y is zero outside of A?

Response: Remark that y has to have Dirac distribution in this case since A is a set of deterministic points, and inclusion y ∈ A has to be interpreted here as “y is an element of A”, and not “the range of y lies in A” with probability 1. This is an important difference, and it is one of the reasons this whole discussion section is
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placed in the paper. Since elements of $A$ are only deterministic points, and they are equivalent to random variables with Dirac probability distributions concentrated at those points, $y$ has to have a Dirac probability distribution if it is an element of $A$.

**Below Eq (20): How do you estimate this with Monte Carlo method? Do you mean some chance constrained optimization?**  
**Response:** The estimate is indeed given in Theorem 10. Indeed, this result may be interpreted as being similar to the chance constrained optimization.  

Intuitively, how fast the system "forgets" the past state depends on its time constant. This is captured by your Lyapunov function. If the dynamics is simple enough, e.g., Ornstein-Uhlenbeck process with asymptotically stable drift, could you comment on how $N$ is affected by dominant eigenvalue of the $A$ matrix of the linear drift? This result can be useful, e.g., for a user to judge (without having to first obtain the Lyapunov function) whether discretizing the space of input history is better than discretizing the state-space.  

**Response:** We thank the reviewer for raising this interesting comment. Unfortunately, a linear stochastic control system may not be $\delta$-ISS-$M_q$, for $q \in \{1, 2\}$, even if matrix $A$ in the drift term is Hurwitz because of the effect of the diffusion term. Instead, one can solve the LMI (4) in the revised version for some positive definite matrix $P$ and largest possible positive constant $\hat{\kappa}$. Here, the constant $\hat{\kappa}$ plays the role of "dominant eigenvalue" to which the reviewer is referring. We believe solving this LMI for linear stochastic control systems is simple and can be quite fast.

Response to Reviewer 2

Essentially, the claim is that you can have a time-discretized stochastic process with the identity as the output map, and a symbolic abstraction with a stochastic output map which essentially replicates the stochasticity of the original dynamics, and you can show that the two systems are approximately bisimilar relative to a $q$-moment metric.  

**Response:** We should point out that we have proposed two types of symbolic abstractions in this paper: one with probabilistic output values and one with deterministic output values. We have shown that both of them are approximately bisimilar to the original stochastic control system.

First, the $q$-moment metric alone is quite weak to capture what happens with the dispersion of the process’ solutions. For example, indeed, you can track the expectations and you may show that the expectation of the stochastic system approaches asymptotically the solution of the associated deterministic system, but over time the process could be all over the place so the expectation convergence could be meaningless. You may then attempt to use the second moment to control dispersion, but then you have no clue what the expectation does.  

**Response:** We thank the reviewer for raising this concern. Although we establish the closeness of a stochastic control system to its symbolic abstraction via the $q$th
moment metric, one can use the results in [31, Propositions 5.9 and 5.11] to infer their closeness in probability as well. Furthermore, we should point out that the closeness in higher moment metric also implies the closeness in the lower moment metric due to Jensen inequality [16]:

\[ E[\|x - x'|^k] \leq \left( E[\|x - x'|^q] \right)^{\frac{k}{q}}, \]

for any \( q \geq k \geq 1 \). In conclusion, what the reviewer has described about the first and second moment metric is not of concern.

A second issue is related to the relaxation of the specification formula done to accommodate randomness. The Euclidean space where the deterministic output may take values in is presented as a (zero measure ?) subset in a set of random variables—I think we are to understand those as functions; so the reference to (non-probabilistic) points is kind of confusing, but this is a different issue that relates to presentation. So set inclusion, as a criterion for specification satisfaction, is relaxed to only something like epsilon-closeness to the target set. Naturally, the issues with the metric are compounded. Another relaxation is encountered in section 5.3 as it relates to state labeling, and stems from the need to relax the formula satisfaction concept. Yet another relaxation is found in the subsequent paragraph where the distance to the specification is to be understood in a randomized (Monte Carlo) way, because obviously there are no closed form solutions for the expectation of the solution of the stochastic process. By this point, it is hard to assess just how much approximation one has introduced in the whole process. Not to mention the inherent conservativeness of the ISS constructions.

Response: We thank the reviewer for the several remarks. We agree with the reviewer in that \( \mathbb{R}^n \)-valued random variables are indeed functions from \( \Omega \) to \( \mathbb{R}^n \). Here, with a slight abuse of notation, we consider any point \( y \in \mathbb{R}^n \) as a random variable with a Dirac probability distribution centred at \( y \), i.e. \( y : \hat{\Omega} \to \{ y \} \) where \( \hat{\Omega} \subset \Omega \) and \( |\hat{\Omega}| = 1 \). We elaborated on this in the revised manuscript. Since requiring set inclusion for deterministic sets and random variables is very conservative (maybe impossible), one is undeniably required to ask for closeness of random variables to deterministic sets justifying the use of approximate bisimulation rather than exact bisimulation. Other relaxations will be applied only for the symbolic models with probabilistic output values. We agree with the reviewer in that the sufficient conditions proposed in (15) and (17) can result in conservative symbolic models depending on the choice of \( \delta\text{-ISS-M}_q \) functions. As we discussed in the paragraph after Theorem 1, one can search for a less conservative \( \delta\text{-ISS-M}_q \) function by resorting to some convex optimization tools, e.g. SOSTOOLS [17] as long as the drift and diffusion term in the system are polynomial functions.

On the presentation side, one may want to be a little more careful in their reference to random variables, points (non-probabilistic / probabilistic) etc. For instance, non-probabilistic does not necessarily mean deterministic. The identification of random variables with points (especially in proximity with references to \( \mathbb{R}^n \) in the text) may be misleading; maybe the -outcome- of a random variable, its image for a given sample point?
Response: We thank the reviewer for the remarks. We changed “non-probabilistic” to “deterministic” in the revised manuscript. We kindly refer the reviewer to the response of the previous remark.

In section 5, I suppose there is some typo in calling \( S(\Sigma) \) deterministic, even in the sense of Definition 4 (which by the way makes no reference to determinism itself - this is later on). System \( S(\Sigma) \) presented in the beginning of section 5.1 is definitely non-deterministic: its transitions involve the solution of a diffusion, so the Post cannot be a singleton. Going back to the Post introduction in section 4, I cannot see how \( |\text{Post}(x)| < 1 \) if this post exists.

Response: We thank the reviewer for the comments. We modified Definition 4 to include the description of determinism as well. Since the set of states in \( S(\Sigma) \) is the set of all \( \mathbb{R}^n \)-valued random variables (each state is a random variable) and the solution process of \( \Sigma \) is uniquely determined (starting from one random variable at \( t = 0 \) results in one random variable at \( t = \tau \)), \( S(\Sigma) \) is deterministic.

In Lemma 1, I cannot (readily) see how \( h_x \) goes to 0 when \( t \to \infty \)

Response: Since \( \beta \in \mathcal{KL} \) in the expression of \( h_x(t) \), the term integral in \( h_x(t) \) is increasing with time \( t \) in at most a linear rate. Therefore, multiplication of an exponentially decreasing function of time (converging to zero) and linearly increasing function of time will converge to zero as time goes to infinity.

The statement in the beginning of section 3.1 is not entirely accurate as stated. This is an upper bound on the -expectation- of the distance

Response: We kindly refer the reviewer to the term “(in the qth moment)” after “the distance” which implies the notion of distance in the expectation.

And before, in Definition 3, I do not think there is any implication in (ii) about the growth rate of \( \bar{\alpha} \) or \( \underline{\alpha} \); just for \( V \)

Response: Since a \( \mathcal{K}_\alpha \) convex function \( \underline{\alpha} \) lower bound a \( \mathcal{K}_\alpha \) concave function \( \bar{\alpha} \), one can infer that the growth rate of those functions should be linear. We point out that \( V : \mathbb{R}^n \times \mathbb{R}^n \to \mathbb{R}_0^+ \) is a function defined over \( \mathbb{R}^n \times \mathbb{R}^n \) and a growth rate is not quite well-defined. Using (ii) in Definition 3 one can infer that \( V(x,x') \) increases proportionally to \( \|x - x'\|^q \) as \( \|x - x'\| \) increases.

Response to Reviewer 3

We have incorporated all the reviewer’s suggestions in the revised version.

Response to Reviewer 4

The notions of stochastic control systems as well as of incremental input-to-state stability in the q-th moment will not be known to a substantial part of the DEDS community. Please explain them by means of examples.
Response: We thank the reviewer for the suggestion: we have elaborated on the notions in the revised manuscript by providing the example of linear stochastic control systems and the corresponding condition ensuring their incremental stability in the first and second moment metric.

The particular notion of input-to-state stability facilitating your approach seems to be a rather strong one. Most readers might deem mere feedforward control devoid of any state observation unreasonable for stochastic systems. Yet this is exactly what your notion of ISS buys: both the dependence on internal state and effects of stochastic disturbances vanish so rapidly that the current state of the controlled process can sufficiently precisely be determined in terms of (reasonably short) recent input sequences only, which in turn facilitates your abstraction. In the light of the fundamental role of this seemingly strong assumption, you ought to discuss the scope of your work, i.e., how likely the assumption is to hold in different domains.

Response: We thank the reviewer for the remark. We point out that our notion
\( \delta \)-ISS-M\(_2\) is nothing more than asymptotic stability in the mean square sense for linear stochastic control systems described in Example 1. For nonlinear stochastic control systems, \( \delta \)-ISS-M\(_q\) is a stronger property than the usual stability. We elaborated on this in the revised version.

The intuition behind the set A-epsilon should be explained in the stochastic case, as there are two other natural (and intrinsically different) possible interpretations of the tolerance parameter epsilon: as a mere metric tolerance in outputs or as a tolerance in probability distributions. Both come together here in the form of an expected tolerance, which however may easily be missed by some readers.

Response: We thank the reviewer for pointing this out. We elaborated on this fact. Note though, that epsilon refers to a distance from the relevant set in the well-defined metric, no matter whether this metric originates from a stochastic system or a deterministic one. The only difference is that in the former case the states are random variables, yet when talking about analysis it only matters that we work over a metric space, and probabilistic nature of the origination of the states is irrelevant.

The model seems somewhat artificial for at least two reasons: First of all, a continuous model based on traffic densities rather than particle (i.e., car) counts is a standard if substantial numbers of cars are involved per segment, yet would generally be assumed to be too coarse a relaxation for small numbers of cars, like the 0-15 range covered here. Second, the underlying notion of ISS, which renders feedforward control w/o state observation possible and thus facilitates the particular abstraction employed in your construction, might be considered an unreasonably strong assumption in traffic control, given that most current traffic light schemes employ forms of traffic and thus state observation for adaptivity, i.e., feedback control. You may want to explain the particular structural properties of your road segment facilitating mere feedforward control.

Response: We thank the reviewer for the remarks. We should point out that the road traffic model we used in this paper is borrowed from [5] and described in details
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in [27] is actually a macroscopic cell transmission model based on the car conservation principle. We kindly refer the reviewer to [27] for more detailed information about the model. Note that since in this specific example the variables $\xi_i$ denote the number of vehicles in cell $i$ with the length of only 0.25 kilometre, the desired range of cars scales to maximum of 15 vehicles. However, one could increase this range by simply increasing the length of each cell to several kilometres. Since the road traffic model is $\delta$-ISS-M$_2$ which is nothing more than asymptotic stability in the mean square sense (cf. 2nd response to the reviewer), the resulting finite abstraction is deterministic and, hence, the feedback implantation and the feedforward implementation of symbolic controllers provide the same behaviours for the road traffic model. One can implement a receding horizon framework on top of the symbolic controller in order to provide a robust solution if the model is not accurate.

I noticed that the safety property is occasionally violated in the sample runs plotted in figure 3. Obviously, there are two different possible reasons for that which interact here, namely that the tolerance in expectations (first cause) is only guaranteed up to a certain confidence (second cause). An explanation of what frequencies of violations are to be expected due to the two reasons might help understanding this fact better. Furthermore, some further evaluation regarding the actual frequency of violation compared to the expected frequency would be worthwhile. What confidence value has been used to synthesize the controller?

Response: We thank the reviewer for the comments. Since we constructed the symbolic model with deterministic output values for the road traffic model, we do not deal with the confidence level at all. Since the symbolic model is $\varepsilon$-approximate bisimilar (in the 2nd moment metric) to the original concrete system, the closed-loop trajectories are always within $\varepsilon$-distance from the safe set justifying the violation at some points in time. Figure 3 (bottom right panel) illustrates the distance of the closed-loop trajectory w.r.t. set $W$ which is always lower than $\varepsilon$. In general, one cannot infer anything regarding the frequency of violation in the context of symbolic control because the inputs are extracted from the evolution of a finite automaton and depend on the initial condition of the plant.

The control input can be either 0 or 6 for $u_1$ and 0 and 8 for $u_2$, why not allowing for arbitrary values between those values. Also, why are the $u$-vectors 5 dimensional, and not 2-dimensional as there are only two control points. The effect of this two-dimensional input would be characterized by a $B \in \mathbb{R}^{2 \times 5}$ matrix. As it is written, 3 of the columns in $B$ seem to have no effect on the dynamics.

Response: Since using these three quantized values we can find a controller satisfying the given specification, one does not need to use finer quantization parameter for the input set. Otherwise, we should choose more quantized values in the intervals $[0, 6]$ and $[0, 8]$ for $u_1$ and $u_2$, respectively. We agree with the reviewer that only a 2 dimensional input vector suffices for this example. We revised matrix $B$ in the revised manuscript.

Is there a special reason choosing the source state different from the initial condition state? It seems to be related to equation 13, but this was not mentioned.
Response: We should point out that in general the source state $x_s$ may not belong to the set of initial conditions in the concrete system. We kindly refer the reviewer to Theorems 6 and 7 for the definition of the set of initial conditions in the concrete systems.

The particular example is illustrative, but given that mere feedforward control may seem rather weak in a stochastic setting, further positive examples as well as modifications demonstrating the limits would be welcome.

Response: The example is aimed mainly at elucidating the technical details of this work and is taken from relevant literature. We also refer the reviewer to the 3rd response in this subsection.

As stated, an input discretization is more desirable whenever the input-space discretization is of smaller cardinality than the state discretization (unnumbered equation page 21, line 25). To me, it is unclear what property of the system will cause this inequality to hold, i.e. is this likely to hold across a large range of systems?

Response: The inequality (22) or the one before it are likely to be satisfied if the state-space dimension of the system is very large, i.e. $n >> 1$, or the system is strongly $\delta$-ISS-$M_\kappa$, i.e. $\kappa >> 1$. We elaborated on this in the revised version.

Although the manuscript is overall well-written, phrases such as 'note that', 'notice that' could be reduced or diversified (currently 29 times across the manuscript).

Response: We revised the manuscript according to the reviewer’s suggestion.

Large parts of the manuscript were already presented in references [26] and [28]. I feel that the additional proofs of the statements of 28 and the new example are not adding significant new insights.

Response: We kindly refer the reviewer to Subsection 5.4 in the revised version for a detailed comparison between the results proposed here and the ones in [31] ([26] in the old version). This paper provides a detailed and extended elaboration of the results announced in [34] ([28] in the old version), including the proofs of the main results, a detailed discussion on how to deal with probabilistic output values and an extension of the corresponding result with no requirement on compactness anymore.

In Section 5.2, the proposed method is only compared to a state-discretization approach presented in 26. Further comparisons to the scaling properties of other approaches (for example the ones mentioned in the introduction) as well as optimal control approaches via Hamilton Jacobi Bellman approaches would make a more significant contribution.

Response: To the best of the authors’ knowledge, only the results proposed in this paper and in [31] ([26] in the old manuscript) provide formal synthesis schemes against general linear temporal logic properties for continuous-time continuous-space stochastic control systems. In order to have a fair and acceptable comparison, we did not compare our results with other existing techniques on stochastic systems.
I could not follow the proof of Theorem 10. Specifically, I don’t understand the application of the Chernoff inequality. Is it a Markov inequality? How do you obtain the ’M’ in the denominator? There seems to be some telescoping going on, but due to the ’q’, I can’t figure out how M and 2q is obtained. It seems when a Chernoff-type inequality is applied, one would obtain an exponential bound.

Response: We agree with the reviewer in that Chernoff-type bounds are often exponential. However, the bound we use in this work can be found e.g. above formula (1) of [4], just above the words “Chernoff’s bounding method is especially convenient ...”. We have added an explicit reference in the new version of the manuscript.

The stated bound in Theorem 10 is different to the one stated in [28], if both are valid: the one in [28] scales better as a function of |A|.

Response: This is indeed true. The previous bound was only derived for a bounded domain, so that the Hoeffding’s inequality could be applied, and only for the case $q = 1$. The new bound in this manuscript is much more general, and hence does not scale as well in the special case considered before.

It seems that an inequality of the type in eq. (33) ((34) in the revised version) can be directly obtained from Hoeffding’s inequality, however, this would require a bound on the support which holds with probability 1, which seems missing in the inequality. If this is indeed not necessary, this would be worth mentioning.

Response: Indeed, as mentioned in the previous response, we could have used this method in [34] ([28] in the old version). Now, as we proceeded for an unbounded case, Hoeffding’s inequality was not applicable anymore. We have pointed this out explicitly in the revised version.