Flower Pollination Algorithm for Global Optimization

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Abstract

Flower pollination is an intriguing process in the natural world. Its evolutionary characteristics can be used to design new optimization algorithms. In this paper, we propose a new algorithm, namely, flower pollination algorithm, inspired by the pollination process of flowers. We first use ten test functions to validate the new algorithm, and compare its performance with genetic algorithms and particle swarm optimization. Our simulation results show the flower algorithm is more efficient than both GA and PSO. We also use the flower algorithm to solve a nonlinear design benchmark, which shows the convergence rate is almost exponential.

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1 Introduction

Nature has been solving challenging problems over millions and billions of years, and many biological systems have evolved with intriguing and surprising efficiency in maximizing their evolutionary objectives such as reproduction. Based on the successfully characteristics of biological systems, many nature-inspired algorithms have been developed over the last few decades [18, 20]. For example, genetic algorithms were based on the Darwinian evolution of biological systems [9] and particle swarm optimization was based on the swarm behaviour of birds and fish [11, 12], which bat algorithm was based on the echolocation behaviour of microbats [21] and firefly algorithm was based on the flashing light patterns of tropic fireflies [18, 19]. All these algorithms have been applied to a wide range of applications.

In many design applications in engineering and industry, we have to try to find the optimal solution to a given problem under highly complex constraints. Such constrained optimization problems are often highly nonlinear, to find the optimal solutions is often a very challenging task if it is not impossible. Most conventional optimization do not work well for problems with nonlinearity and multimodality. Current trend is to use nature-inspired metaheuristic algorithms to tackle such difficult problems, and it has been shown that metaheuristics are surprisingly very efficient. For this reason, the literature of metaheuristics has expanded tremendously in the last two decades [18, 20]. Up to now, researchers have only use a very limited characateristics inspired by nature, and there is room for more algorithm development.

In this paper, we will propose a new algorithm based on the flower pollination process of flowering plants. We will first briefly review the main characteristics of flower pollination,
and thus idealize these characteristics into four rules. We will then use them to develop a flower pollination algorithm (FPA), or the flower algorithm. Then, we validate it using a set of well-known test functions and design benchmark. We analyze the simulations and make comparison of its performance with genetic algorithm and particle swarm optimization. Finally, we discuss further topics for extending this algorithm.

From the biological evolution point of view, the objective of the flower pollination is the survival of the fittest and the optimal reproduction of plants in terms of numbers as well as most fittest. This is in fact an optimization process of plant species. All the above factors and processes of flower pollination interact so as to achieve optimal reproduction of the flowering plants. Therefore, this can inspire to design new optimization algorithm. The basic idea of flower pollination in the context of bees and clustering was investigated before [10], but in this paper, we will design a completely new optimization solely based on the flower pollination characteristics.

2 Characteristics of Flower Pollination

It is estimated that there are over a quarter of a million types of flowering plants in Nature and that about 80% of all plant species are flowering species. It still remains partly a mystery how flowering plants came to dominate the landscape from Cretaceous period [16, 22]. Flowering plant has been evolving for more than 125 million years and flowers have become so influential in evolution, we cannot image how the plant world would be without flowers. The main purpose of a flower is ultimately reproduction via pollination. Flower pollination is typically associated with the transfer of pollen, and such transfer is often linked with pollinators such as insects, birds, bats and other animals. In fact, some flowers and insects have co-evolved into a very specialized flower-pollinator partnership. For example, some flowers can only attract and can only depend on a specific species of insects for successful pollination [7].

Pollination can take two major forms: abiotic and biotic. About 90% of flowering plants belong to biotic pollination, that is, pollen is transferred by a pollinator such as insects and animals. About 10% of pollination takes abiotic form which does not require any pollinators. Wind and diffusion in water help pollination of such flowering plants and grass is a good example [14, 7]. Pollinators, or sometimes called pollen vectors, can be very diverse. It is estimate there are at least 200,000 variety of pollinators such as insects, bats and birds.

Honeybees are a good example of pollinator, and they can also developed the so-called flower constancy [3]. That is, these pollinators tend to visit exclusive certain flower species while bypassing other flower species. Such flower constancy may have evolutionary advantages because this will maximize the transfer of flower pollen to the same or conspecific plants, and thus maximizing the reproduction of the same flower species. Such flower constancy may be advantageous for pollinators as well, because they can be sure that nectar supply is available with their limited memory and minimum cost of learning or exploring. Rather than focusing on some unpredictable but potentially more rewarding new flower species, flower constancy may require minimum investment cost and more likely guaranteed intake of nectar [17].

Pollination can be achieved by self-pollination or cross-pollination. Cross-pollination, or allogamy, means pollination can occur from pollen of a flower of a different plant, while self-pollination is the fertilization of one flower, such as peach flowers, from pollen of the same flower or different flowers of the same plant, which often occurs when there is no
reliable pollinator available.

Biotic, cross-pollination may occur at long distance, and the pollinators such as bees, bats, birds and flies can fly a long distance, thus they can considered as the global pollination. In addition, bees and birds may behave as Lévy flight behaviour [13], with jump or fly distance steps obey a Lévy distribution. Furthermore, flower constancy can be used an increment step using the similarity or difference of two flowers.

3 Flower Pollination Algorithm

Now we can idealize the above characteristics of pollination process, flower constancy and pollinator behaviour as the following rules:

1. Biotic and cross-pollination is considered as global pollination process with pollen-carrying pollinators performing Lévy flights.

2. Abiotic and self-pollination are considered as local pollination.

3. Flower constancy can be considered as the reproduction probability is proportional to the similarity of two flowers involved.

4. Local pollination and global pollination is controlled by a switch probability $p \in [0, 1]$. Due to the physical proximity and other factors such as wind, local pollination can have a significant fraction $p$ in the overall pollination activities.

Obviously, in reality, each plant can have multiple flowers, and each flower patch often release millions and even billions of pollen gametes. However, for simplicity, we also assume that each plant only has one flower, and each flower only produce one pollen gamete. Thus, there is no need to distinguish a pollen gamete, a flower, a plant or solution to a problem. This simplicity means a solution $x_i$ is equivalent to a flower and/or a pollen gamete. In future studies, we can easily extend to multiple pollen gametes for each flower and multiple flowers for multiobjective optimization problems.

From the above discussions and the idealized characteristics, we can design a flower-based on algorithm, namely, flower pollination algorithm (FPA). There are two key steps in this algorithm, they are global pollination and local pollination.

In the global pollination step, flower pollens are carried by pollinators such as insects, and pollens can travel over a long distance because insects can often fly and move in a much longer range. This ensures the pollination and reproduction of the most fittest, and thus we represent the most fittest as $g_*$. The first rule plus flower constancy can be represented mathematically as

$$x_{i}^{t+1} = x_{i}^{t} + L(x_{i}^{t} - g_*),$$

where $x_{i}^{t}$ is the pollen $i$ or solution vector $x_i$ at iteration $t$, and $g_*$ is the current best solution found among all solutions at the current generation/iteration. The parameter $L$ is the strength of the pollination, which essentially is a step size. Since insects may move over a long distance with various distance steps, we can use a Lévy flight to mimic this characteristic efficiently [13, 15]. That is, we draw $L > 0$ from a Levy distribution

$$L \sim \frac{\lambda \Gamma(\lambda) \sin(\pi \lambda/2)}{\pi} \frac{1}{s^{1+\lambda}}, \quad (s \gg s_0 > 0).$$

(2)
Flower Pollination Algorithm (or simply Flower Algorithm)

**Objective**

$$\text{min or max } f(x), \ x = (x_1, x_2, ..., x_d)$$

*Initialize a population of n flowers/pollen gametes with random solutions*

*Find the best solution $g_*$ in the initial population*

*Define a switch probability $p \in [0, 1]$*

**while** ($t < \text{MaxGeneration}$)

**for** $i = 1 : n$ (all $n$ flowers in the population)

*if* rand $< p$,

*Draw a (d-dimensional) step vector $L$ which obeys a Lévy distribution*

*Global pollination via* $x_i^{t+1} = x_i^t + L(g_* - x_i^t)$

*else*

*Draw $\epsilon$ from a uniform distribution in $[0,1]$*

*Randomly choose $j$ and $k$ among all the solutions*

*Do local pollination via* $x_i^{t+1} = x_i^t + \epsilon(x_j^t - x_k^t)$

*end if*

**Evaluate new solutions**

*If new solutions are better, update them in the population*

**end for**

*Find the current best solution $g_*$*

**end while**

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Figure 1: Pseudo code of the proposed Flower Pollination Algorithm (FPA).

Here $\Gamma(\lambda)$ is the standard gamma function, and this distribution is valid for large steps $s > 0$. In all our simulations below, we have used $\lambda = 1.5$.

The local pollination (Rule 2) and flower constancy can be represented as

$$x_i^{t+1} = x_i^t + \epsilon(x_j^t - x_k^t),$$

where $x_j^t$ and $x_k^t$ are pollens from the different flowers of the same plant species. This essentially mimic the flower constancy in a limited neighborhood. Mathematically, if $x_j^t$ and $x_k^t$ comes from the same species or selected from the same population, this become a local random walk if we draw $\epsilon$ from a uniform distribution in $[0,1]$.

Most flower pollination activities can occur at both local and global scale. In practice, adjacent flower patches or flowers in the not-so-far-away neighborhood are more likely to be pollinated by local flower pollens that those far away. For this, we use a switch probability (Rule 4) or proximity probability $p$ to switch between common global pollination to intensive local pollination. To start with, we can use $p = 0.5$ as an initially value and then do a parametric study to find the most appropriate parameter range. From our simulations, we found that $p = 0.8$ works better for most applications.

The above two key steps plus the switch condition can be summarized in the pseudo code shown in Fig. 1.

4 Numerical Results

Any new optimization should be extensively validated and comparison with other algorithms. There are many test functions, at least over a hundred well-know test functions. However, there is no agreed set of test functions for validating new algorithms, though there some review and literature [1, 5, 19]. In this paper, we will choose a diverse subset of such test functions to validate our proposed Flower Pollination Algorithm (FPA).
In addition, we will also compare the performance of our algorithm with that of genetic algorithms [8] and particle swarm optimization [11, 12]. Furthermore, we will also apply FPA to solve a well-known pressure vessel design benchmark [2, 6].

4.1 Test Functions

The Ackley function can be written as

\[ f(x) = -20 \exp \left( -\frac{1}{5} \sqrt{\frac{1}{d} \sum_{i=1}^{d} x_i^2} \right) - \exp \left( \frac{1}{d} \sum_{i=1}^{d} \cos(2\pi x_i) \right) + 20 + e, \]  

which has a global minimum \( f_* = 0 \) at \((0, 0, ..., 0)\).

The simplest of De Jong’s functions is the so-called sphere function

\[ f(x) = \sum_{i=1}^{n} x_i^2, \quad -5.12 \leq x_i \leq 5.12, \]  

whose global minimum is obviously \( f_* = 0 \) at \((0, 0, ..., 0)\). This function is unimodal and convex.

Easom’s function

\[ f(x) = -\cos(x) \cos(y) \exp \left( - (x - \pi)^2 + (y - \pi)^2 \right), \]  

whose global minimum is \( f_* = -1 \) at \( x_* = (\pi, \pi) \) within \(-100 \leq x, y \leq 100\). It has many local minima.

Griewangk’s function

\[ f(x) = \frac{1}{4000} \sum_{i=1}^{n} x_i^2 - \prod_{i=1}^{n} \cos \left( \frac{x_i}{\sqrt{i}} \right) + 1, \quad -600 \leq x_i \leq 600, \]  

whose global minimum is \( f_* = 0 \) at \( x_* = (0, 0, ..., 0)\). This function is highly multimodal.

Michaelwicz’s function

\[ f(x) = -\sum_{i=1}^{n} \sin(x_i) \cdot \left[ \sin \left( \frac{ix_i^2}{\pi} \right) \right]^{2m}, \]  

where \( m = 10 \), and \( 0 \leq x_i \leq \pi \) for \( i = 1, 2, ..., n \). In 2D case, we have

\[ f(x, y) = -\sin(x) \sin^{20} \left( \frac{x^2}{\pi} \right) - \sin(y) \sin^{20} \left( \frac{2y^2}{\pi} \right), \]  

where \((x, y) \in [0, 5] \times [0, 5]\). This function has a global minimum \( f_* \approx -1.8013 \) at \( x_* = (x_*, y_*) = (2.20319, 1.57049)\).

Rastrigin’s function

\[ f(x) = 10n + \sum_{i=1}^{n} \left[ x_i^2 - 10 \cos(2\pi x_i) \right], \quad -5.12 \leq x_i \leq 5.12, \]  

whose global minimum is \( f_* = 0 \) at \((0, 0, ..., 0)\). This function is highly multimodal.
Rosenbrock’s function

\[
f(x) = \sum_{i=1}^{n-1} [(x_i - 1)^2 + 100(x_{i+1} - x_i^2)^2],
\]  

whose global minimum \( f_* = 0 \) occurs at \( x_* = (1, 1, ..., 1) \) in the domain \(-5 \leq x_i \leq 5\) where \( i = 1, 2, ..., n \). In the 2D case, it is often written as

\[
f(x, y) = (x - 1)^2 + 100(y - x^2)^2,
\]

which is often referred to as the banana function.

Schwefel’s function

\[
f(x) = -\sum_{i=1}^{n} x_i \sin(\sqrt{|x_i|}), \quad -500 \leq x_i \leq 500,
\]

whose global minimum \( f_* \approx -418.9829n \) occurs at \( x_i = 420.9687 \) where \( i = 1, 2, ..., n \).

Yang’s forest-like function \([20]\)

\[
f(x) = \left( \sum_{i=1}^{d} |x_i| \right) \exp \left[ -\sum_{i=1}^{d} \sin(x_i^2) \right], \quad -2\pi \leq x_i \leq 2\pi,
\]

has a global minimum \( f_* = 0 \) at \( (0, 0, ..., 0) \).

Shubert’s function

\[
f(x) = \left[ \sum_{i=1}^{n} i \cos \left( i + (i+1)x \right) \right] \cdot \left[ \sum_{i=1}^{n} i \cos \left( i + (i+1)y \right) \right],
\]

which has 18 global minima \( f_* \approx -186.7309 \) for \( n = 5 \) in the search domain \(-10 \leq x, y \leq 10\).

Table 1: Comparison of algorithm performance in terms of number of iterations.

| Functions/Algorithms | GA     | PSO    | FPA    |
|----------------------|--------|--------|--------|
| Michalewicz (d=16)  | 89325 ± 7914(95%) | 6922 ± 537(98%) | 3341 ± 649(100%) |
| Rosenbrock (d=16)   | 55723 ± 8901(90%) | 32756 ± 5325(98%) | 5532 ± 1464(100%) |
| De Jong (d=256)     | 25412 ± 1237(100%) | 17040 ± 1123(100%) | 4245 ± 545(100%) |
| Schwefel (d=128)    | 227329 ± 7572(95%) | 14522 ± 1275(97%) | 6851 ± 448(100%) |
| Ackley (d=128)      | 32720 ± 3327(90%) | 23407 ± 4325(92%) | 3357 ± 968(100%) |
| Rastrigin           | 110523 ± 5199(77%) | 79491 ± 3715(90%) | 10840 ± 2689(100%) |
| Easom               | 19239 ± 3307(92%) | 17273 ± 2929(90%) | 4017 ± 982(100%) |
| Griewank            | 70925 ± 7652(90%) | 55970 ± 4223(92%) | 4918 ± 1429(100%) |
| Yang (d = 16)       | 27923 ± 3025(83%) | 14116 ± 2949(90%) | 4254 ± 1839(100%) |
| Shubert(18 minima)  | 54077 ± 4997(89%) | 23992 ± 3755(92%) | 9271 ± 1758(100%) |

For the above ten test functions, each function can have varied dimensions; then there is an issue which dimensions should be used in the simulations. Research suggests that higher-dimensional problems tend to be more challenging, and a new algorithm should be
tested against a wide range of functions in terms of function properties and dimensions. Therefore, we tend to focus on problems with higher dimensions.

In addition, we have used three algorithms to find their optimal solution with a given tolerance $10^{-5}$. The three algorithms are genetic algorithm (GA), particle swarm optimization (PSO) and the new flower pollination algorithm (FPA). For each algorithm, we have carried out 100 independent runs using a population size $n = 25$ and $p = 0.8$ for FPA, crossover probability 0.95 and mutation probability 0.05 for GA, and learning parameters 2 for PSO. The results are summarized in Table 1. In the table, the results are provided as mean ± standard deviation (success rate). For example, $3341 ± 649 (100\%)$ means that mean number iterations is $3341$ with one standard deviation of $649$ and a success rate of 100%. The total number of function evaluations is $n$ times the mean number of iterations. For example, the number of iterations is $3341$ in the table, so the total number of function evaluations is $3341 \times 25 = 83525$.

4.2 Design Optimization

Pressure vessels are literally everywhere such as champagne bottles and gas tanks. For a given volume and working pressure, the basic aim of designing a cylindrical vessel is to minimize the total cost. Typically, the design variables are the thickness $d_1$ of the head, the thickness $d_2$ of the body, the inner radius $r$, and the length $L$ of the cylindrical section [2]. This is a well-known test problem for optimization and it can be written as

$$\text{minimize } f(x) = 0.6224d_1rL + 1.7781d_2r^2 + 3.1661d_1^2L + 19.84d_2^2r,$$  \hspace{1cm} (16)

subject to the following constraints

$$g_1(x) = -d_1 + 0.0193r \leq 0$$
$$g_2(x) = -d_2 + 0.00954r \leq 0$$
$$g_3(x) = -\pi r^2L - \frac{4\pi}{3}r^3 + 1296000 \leq 0$$
$$g_4(x) = L - 240 \leq 0.$$  \hspace{1cm} (17)

The simple bounds are

$$0.0625 \leq d_1, d_2 \leq 99 \times 0.0625,$$  \hspace{1cm} (18)

and

$$10.0 \leq r, \quad L \leq 200.0.$$  \hspace{1cm} (19)

Recently, Cagnina et al (2008) used an efficient particle swarm optimiser to solve this problem and they found the best solution

$$f_* \approx 6059.714,$$  \hspace{1cm} (20)

at

$$x_* \approx (0.8125, 0.4375, 42.0984, 176.6366).$$  \hspace{1cm} (21)

This means the lowest price is about $6059.71.

Using the proposed flower pollination algorithm, we have easily found the same as the solution $f_* \approx 6059.714$ obtained by Cagnina et al [2, 6].

The current best solution can be stored during iterations. Then, we can calculate the errors $D$ in terms of the difference between the current solution to the best mean solution.
after 40 independent runs. Figure 2 shows the typical variations of \( D \) during iterations. We can also see that the proposed algorithm approaches the optimal solution exponentially (see Fig. 2).

Among the three methods, the proposed FPA obtained the best result and converged most quickly.

5 Discussions

Flowering plants have evolved some interesting features of flower pollination, and we have successfully developed a new flower algorithm to mimic these characteristics. Our simulation results have shown that the the proposed flower pollination algorithm is very efficient and can outperform both genetic algorithm and particle swarm optimization. The convergence rate is essentially exponential as we have seen from the convergence comparison in the previous section.

The reasons that FPA is efficient can be twofold: long-distance pollinators and flower consistency. Pollinators such as insects can travel long distance, and thus they introduce the ability (into the algorithm) that they can escape any local landscape and subsequently explore larger search space. This acts as exploration moves. On the other hand, flower consistency ensure that the same species of the flowers (thus similar solutions) are chosen more frequently and thus guarantee the convergence more quickly. This step is essentially an exploitation step. The interplay and interaction of these key components and the selection of the best solution \( g^* \) ensure that the algorithm is very efficient.

6 Conclusions

In the present algorithm, for simplicity, we have assumed that each flower only produce one pollen gamete, this simplifies the implementation greatly. However, to assign each flower with multiple pollen gametes and each plant with multiple flowers can have some advantages for some applications such as image compression, multiobjective optimization, and graph colouring. This can form a topic for further research.
Another possible extension is to use design other schemes for flower constancy. At present, this constancy is realized by a simple formula. Other more exotic form may be useful to certain type of problem, though the exact improvement may need some extensive simulations.

Furthermore, it is possible to extend the flower algorithm to a discrete version so that it can solve combinatorial optimization problems. All these extensions will be very useful. We hope that this paper will inspire more active research in metaheuristics in the near future.

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