Sufficient Conditions and Constraints for Reversing General Quantum Errors

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Reversing the effects of a quantum evolution, for example as is done in error correction, is an important task for controlling quantum systems in order to produce reliable quantum devices. When the evolution is governed by a completely positive map, there exist reversibility conditions, known as the quantum error correcting code conditions, which are necessary and sufficient conditions for the reversibility of a quantum operation on a subspace, the code space. However, if we suppose that the evolution is not described by a completely positive map, necessary and sufficient conditions are not known. Here we consider evolutions that do not necessarily correspond to a completely positive map. We prove the completely positive map error correction conditions can lead to a code space that is not in the domain of the map, meaning that the output of the map is not positive. A corollary to our theorem provides a class of relevant examples. Finally, we provide a set of sufficient conditions that will enable the use of quantum error correcting code conditions while ensuring positivity.

I. INTRODUCTION

Reversing quantum operations is an important form of quantum control, which will help to enable many quantum technologies. For example, error correction, which is the reversal of an unwanted quantum operation, will be necessary to ensure that errors do not ruin a quantum computer’s algorithm execution. Error correction is also important in long distance communication to ensure data integrity. Quantum error correction was shown to be possible, and potentially practical, with the invention of the Shor [1] and Steane [2] quantum error correcting (QEC) codes. Subsequently, with a set of reasonable assumptions, necessary and sufficient conditions for the existence of an error correcting code were provided by Bennett et. al. [3], Knill and Laflamme [4], and Nielsen et al. [5]. These conditions may, more generally, be seen as reversibility conditions since they indicate when it is possible to reverse the effects of an evolution.

Such conditions are often described in terms of a completely positive (CP) map, $A$. That is, a map that maps all positive operators to positive operators and does so even when extended by an identity operator to $I \otimes A$. This is sometimes also called a dynamical map although not all maps are completely positive (e.g. the transpose) and the terminology is not consistent in the literature with respect to dynamical maps. It should also be noted that there is an ongoing discussion in the Physics community about the physicality of non-completely positive (NCP) maps. (See for example [6] and references therein.) However, most researchers consider a map physical if the domain of the map is restricted to positive output density operators $\rho \rightarrow \rho'$. Without directly addressing this problem here, we present conditions which restrict the ability to reverse an evolution that does not correspond to a CP map. In other words, the evolution we consider is not required to correspond to a CP map. Furthermore, it may not correspond to any map at all, but rather to a single system evolving in time from one state to another.

When we study the reversibility of a system, particularly for error correction, we are often looking at a subspace $\mathcal{H}_S$ of the system-environment Hilbert space $\mathcal{H}_{SE}$, where $E$ is the environment and $S$ the system. The initial state of the system is $\rho = \operatorname{tr}_E(\rho_{SE})$, where $\rho_{SE}$ is the initial combined system and environment state. One can experimentally determine a dynamical map that describes the open-system evolution of the system under consideration $A : \rho_S \rightarrow \rho'_S$. This will determine the set of errors that occur on the system and an appropriate error-correcting code can be determined from the set of errors that are targeted for correction. It is well known that when the initial state of the system and environment together is a product state, that is when they are uncorrelated, $\rho_{SE} \equiv \rho_S \otimes \rho_E$ and the evolution can be described by a completely positive map.

The model of error correction that we consider is where the recovery operation is implemented after the error. This is the model usually considered and is, for example, discussed in some detail in the book by Nielsen and Chuang [9]. To be more specific, the process for quantum error correction occurs in four main steps. In the first step, the system is encoded. Next, the system evolves, possibly incurring an error. Then, a measurement is made to extract the error syndrome to identify a possible correctable error. Finally, the error, if present, is corrected using a unitary transformation. The way to express this, arising from Eq. (9) (below), is $U_k P_k$ where $P_k$ is the measurement to detect an error and $U_k$ is the corresponding unitary which is implemented conditioned on the outcome of the measurement $P_k$. A diagram of an example of this process for the single bit-flip repetition code is shown in Figure 1. The details of the gates in Figure 1 are not important for our situation. However, it should be emphasized that the correction process depends on the syndrome measurement outcome.

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When the evolution does not correspond to a CP map, we would like to find a way to generalize or extend the reversibility conditions (CP error correcting conditions [9]) to more general evolutions. Motivated by a desire to describe very general error models, such as those consid-
ered by Aharanov and Ben-Or [10], Shabani and Lidar [11] studied this problem and showed that the same code
space for the corresponding CP map works for a corre-
sponding NCP map (specified below), but they focused on the hermiticity of the evolution and not the positiv-
ity. In this paper, we show that if an evolution is not described by a CP map, satisfying the CP quantum error correcting conditions can produce a code space that is not in the domain of the NCP error map in the sense that it does not produce a positive output. In contrast to [11], we seek an output that is not only hermitian, but also positive. We provide conditions on the code, via Theorem 1, such that the quantum error correcting conditions for an NCP evolution will produce a positive, hermitian output. This leads to a set of sufficient conditions for the reversibility of an NCP evolution when we demand that the output be both hermitian and positive. This is our Theorem 2 which is followed by examples.

II. BACKGROUND

A superoperator \( \mathcal{A} \) can be represented by a matrix acting on \( \rho_S \) [12]

\[
\rho_{s',s} = \mathcal{A}^\dagger_{s',s} \rho_{s}. \tag{1}
\]

(The sum over repeated indices is implied.) The evolu-
tions we consider will be those that preserve the hermitic-
ity and trace. In this case the matrix \( \mathcal{A} \) must satisfy the conditions, respectively,

\[
\mathcal{A}^{\dagger}_{s',s} = (\mathcal{A}_{s',s})^* \tag{2}
\]

and

\[
\mathcal{A}_{s',s'} = \delta_{s,s}. \tag{3}
\]

For an alternative description, we often use the matrix \( \mathcal{B} \), which is related to the \( \mathcal{A} \) matrix by

\[
\mathcal{B}_{s',s',s} = \mathcal{A}_{s',s'} \rho_s. \tag{4}
\]

The hermiticity condition (2) translates to

\[
\mathcal{B}_{s',s',s} = (\mathcal{B}_{s',s'})^*. \tag{5}
\]

Then, a general Hermitian preserving linear map can be written in an operator-sum decomposition \( \mathcal{E} \) of the form

\[
\mathcal{E}(\rho) = \sum_i \eta_i E_i \rho E_i^\dagger, \tag{6}
\]

where the \( \{\eta_i\} \)'s are the signs of the eigenvalues and the \( \{E_i\} \)'s are the eigenvectors of the matrix \( \mathcal{B} \) after absorb-
ing the magnitudes of the eigenvalues [2, 12, 13]. These can be written in matrix form.

If the system is not correlated with the environment, i.e., the combined system and environment is a product state \( \rho_S \otimes \rho_E \), then the evolution of the system is given by a completely positive (CP) map, and all the \( \eta_i \) = 1. In the case that the system and environment are not initially in a product state, general conditions for complete positivity are not known, but in some special cases the map is still CP [13, 21]. However, whenever the map \( \mathcal{E}(\rho) \) is CP, we can write it as \[13, 22\]

\[
\mathcal{E}(\rho) = \sum_i E_i \rho E_i^\dagger. \tag{7}
\]

Furthermore, when the evolution corresponds to a CP map, there is a set of quantum error correcting code con-
ditions, which ensure the reversibility of the evolution. These are necessary and sufficient for the construction of a quantum error correcting code, which can be used to detect and correct the errors, thus reversing the effects of the map. One way of expressing these conditions is \[4\]

\[
(\alpha_L | E_i^\dagger E_j | \beta_L) = m_{ij} \delta_{\alpha \beta}, \tag{8}
\]

where \( |\alpha_L\rangle \) and \( |\beta_L\rangle \) are logical (encoded states) and \( m_{ij} \) is a constant. This equation is easy to interpret. If \( |\alpha_L\rangle \) is acted on by an error \( E_i \) and \( |\beta_L\rangle \) is another state acted on by an error \( E_j \), then the overlap between these must be zero if the states are different. This ensures that a measurement performed to identify the error will not result in an ambiguous correction procedure to recover the original state. This, and other manifestations in classical error correction, are sometimes called the “disjointness condition” since it shows that the subspace of a logical state acted upon by any correctable error must be disjoint, as a set, from any other logical state with a corre-
tectable error acting on it. It is easy to show that these conditions are satisfied if and only if the equivalent neces-
sary and sufficient conditions for error correction for CP maps are satisfied \[2\]

\[
P E_i^\dagger E_j P = c_{ij} P, \tag{9}
\]

where \( P \) is the projector onto the code space and \( c_{ij} \) are elements of a Hermitian matrix.

A system can often develop correlations with its envi-
ronment so that the combined system-environment state
is no longer a product state \( \rho_{SE} \neq \rho_S \otimes \rho_E \). Correlations between the system and environment can be prevented with dynamical decoupling, but dynamical decoupling does not remove correlations that are present prior to the decoupling operations \([23, 24]\). Given a correlated system and environment, the evolution of the system is not necessarily given by a CP map \([25, 26, 27, 28]\). A not completely positive evolution can be described by a \( \mathcal{B} \) matrix that has at least one negative eigenvalue and has the operator-sum decomposition form

\[
\mathcal{E}(\rho) = \sum_i \eta_i E_i \rho E_i^\dagger
\]  

(10)

where the \( \eta_i \)'s are not all positive \([7]\). Such an evolution may not correspond to a map, but only a specific input and output state.

III. REVERSIBILITY CONDITIONS

Our first main theorem shows that we need to be careful when extending results from CP maps to NCP evolution if we want to ensure positivity.

First, let us define a pseudounitary (PU) transformation with signature \( p, q \) to be a matrix \( U \) such that \( U \eta U^\dagger = \eta \), and a pseudohermitian (PH) matrix is given by \( H^\dagger = \eta H \eta^{-1} \), where, in our case, \( \eta = \text{diag}(1, 1, \ldots, 1, -1, -1, \ldots, -1) \) (\( p \) ones and \( q \) negative ones). (For a more general and thorough discussion, see \([28]\) and the Appendix \([\text{A}]\).) There exists a pseudounitary degree of freedom in the operator-sum decomposition that can be used to express the map in terms of a different set of operators as shown in \([30]\). (For completeness, we provide a slightly different proof in Appendix \([B]\) that we believe is clearer.) The first task is to show that the matrix \( c_{ij} \) in Lemma \([1]\) can be diagonalized by choosing a pseudounitary transformation which will transform the \( E_i \) to a new set that produces the same map, but has \( c_{ij} \) diagonal. We call Eq. \((11)\) (below) the pseudohermitian form of the CP error correcting conditions because when you diagonalize it, you get the diagonalized CP error correcting conditions.

**Lemma 1.** Given an NCP map \( \mathcal{E}(\rho) = \sum_i \eta_i E_i \rho E_i^\dagger \), the PH form of the CP error correcting condition

\[ \eta_i P E_i^\dagger E_j P = c_{ij} P, \]

(11)

where \( c_{ij} \) are elements of a pseudohermitian matrix \( C \), can be diagonalized using the pseudounitary degree of freedom, and it leads to the diagonalized CP error correcting condition.

**Proof.** We can choose a PU transformation \( U \) with elements \( u_{kj} \) such that \( F_j = E_k u_{kj} \). In other words, the \( F_j \) are linear combinations of the \( E_k \) with the set of coefficients \( u_{kj} \) forming a PU matrix.

We can diagonalize \( c_{ij} \) in \((11)\) by using the pseudo unitary degree of freedom of the operators. We switch to block matrix notation by letting

\[
F = \begin{bmatrix} F_1 & F_2 & \cdots & F_n \end{bmatrix},
\]

(12)

\[
E = \begin{bmatrix} E_1 & E_2 & \cdots & E_n \end{bmatrix},
\]

(13)

\[
P = I_n \otimes P.
\]

(14)

We can make the number of elements equal in \( F \) and \( E \), by inserting zero matrices. We treat the block components in \((12), (13), \) and \((14)\) as elements so \( F = EU \) and \( \eta P = P \eta \). Letting \( M = E^\dagger E \), we have

\[
\eta P F^\dagger FP = \eta P U^\dagger M U P
\]

\[
= P \eta U^\dagger \eta M U P
\]

\[
= P \eta^{-1} M U P
\]

\[
= U^{-1} \eta P M U \eta
\]

\[
= U^{-1} C U P
\]

\[
= D P,
\]

(15)

where \( D \) is a diagonal pseudohermitian and we used the property that a pseudounitary matrix can diagonalize a pseudohermitian (see the Appendix \([A]\) and \([B]\) for details). Since \( D \) is diagonal, we have \( D = \eta D^\dagger \eta = D^\dagger \eta = D^\dagger \). Thus, \( D \) is also Hermitian. Then, we can bring \( \eta \) to the right hand side and absorb it into \( D \) because \( \eta D \) is also a diagonal Hermitian. We can simply write

\[
PF^\dagger FP = DP
\]

(16)

or in index notation

\[
PF_j^\dagger F_j P = d_{ij} \delta_{ij} P,
\]

(17)

where \( D \) is a diagonal Hermitian. This is the same as the diagonalized quantum error correction condition for CP maps.

Note that the ability to diagonalize this matrix is tantamount to finding a set of orthogonal projectors that can be used to define a syndrome measurement.

We can now prove the main theorem.

**Theorem 1.** Let us consider a \( \mathcal{B} \) matrix with at least one negative eigenvalue. Let its action correspond to an NCP map

\[
\mathcal{E}(\rho) = \sum_i \eta_i E_i \rho E_i^\dagger,
\]

(18)

where not all of the \( \eta_i \)'s are positive. Now suppose \( \exists U \in \text{PU} \) relating two sets of operators \( \{E_i\} \) and \( \{F_j\} \) such that \( \mathcal{E} \) is equivalent to

\[
\mathcal{E}(\rho) = \sum_j \eta_j F_j \rho F_j^\dagger = \mathcal{E}_1(\rho) - \mathcal{E}_2(\rho),
\]

(19)

where \( \mathcal{E}_1 \) and \( \mathcal{E}_2 \) are both CP maps. In addition, assume that \( PF_j^\dagger F_j P = P d_{ij} \delta_{ij} \), i.e., we satisfy the diagonalized error correcting condition. If \( \mathcal{E}_2(P \rho P) \neq 0 \), then the code space is not in the domain of the error map.
Following Nielsen et al. \[5\], we can show that enforcing the diagonal error correcting condition leads to a code space that is not in the domain of the error map. Let our input density matrix be in the code space, which can be written as \(P \rho P\).

Starting from \(\mathbf{19}\), we can use the polar decomposition to get \(F_k P = U_k \sqrt{PF_k \dagger F_k P} = \sqrt{d_{kk}} U_k P\). Therefore, \(F_k\) rotates the code subspace into the subspace given by the projector

\[
P_k \equiv U_k P U_k \dagger = F_k P U_k \dagger / \sqrt{d_{kk}}.
\]

Then

\[
F_k P = \sqrt{d_{kk}} P_k U_k.
\]

The diagonal error correcting condition ensures that these rotated subspaces are orthogonal, since when \(k \neq l\),

\[
P_l P_k = P_l \dagger P_k = \frac{U_l \rho F_l \dagger F_k P U_k \dagger}{\sqrt{d_{ll} d_{kk}}} = 0.
\]

Using \(\mathbf{19}, \mathbf{21}\), and \(\mathbf{22}\), we can measure the output state and we have outcome \(j\) (unnormalized)

\[
\mathcal{M}_j(\rho') = \mathcal{M}_j(\mathcal{E}(P \rho P)) = P_j \mathcal{E}_1(P \rho P) P_j - P_j \mathcal{E}_2(P \rho P) P_j
\]

\[
= \sum_{k \neq l} |d_{kk}| P_j P_k U_k \rho U_k \dagger P_k P_j
\]

\[
- \sum_{k \neq l} |d_{kk}| P_j P_l U_l \rho U_l \dagger P_l P_j.
\]

Since \(\mathcal{E}_2(P \rho P) \neq 0\), using the orthogonality of \(P_i\) we can choose \(j\) to be one of the \(l\) thus \(P_j P_l = P_l\) for one \(l\) and the other \(P_i\) terms vanish. Thus, the probability of the outcome is a negative value. Since for any valid positive semi-definite density operator this is not possible, the code space cannot be in the domain of the error map. \(\square\)

Theorem \(\mathbf{1}\) leads to a useful corollary. In Eq. \(\mathbf{28}\), we relied on the fact that \(\mathcal{E}_2(P \rho P) \neq 0\). Thus, the negative terms are nonzero. It follows that if the \(F\) operators are unitary, as is the case with the Pauli matrices, we also arrive at the restriction as shown in Corollary \(\mathbf{1}\)

**Corollary 1.** For a \(B\) matrix with at least one negative eigenvalue, if the pseudo unitary degree of freedom leads to

\[
PF_j \dagger F_j P = P \delta_{ji},
\]

where the operators \(F_i\) are unitary and \(P\) is the projector onto the code space, then the code space is not in the domain of the error map.

**Proof.** The proof is similar to the proof for Theorem \(\mathbf{1}\).

Let

\[
P_k \equiv F_k P F_k \dagger.
\]

Then

\[
F_k P = P_k F_k.
\]

Equation \(\mathbf{25}\) ensures that these rotated subspaces are orthogonal, since when \(k \neq l\),

\[
P_l P_k = F_l P F_k \dagger P_k F_l \dagger = 0.
\]

For states \(P \rho P\) in the code space, we have \(\mathcal{E}(P \rho P) = \sum_j \eta_j F_j P \rho P F_j \dagger = \mathcal{E}_1(P \rho P) - \mathcal{E}_2(P \rho P)\). Using \(\mathbf{27}\) and \(\mathbf{28}\), we measure the state and we have outcome \(j\) (unnormalized)

\[
\mathcal{M}_j(\rho') = \mathcal{M}_j(\mathcal{E}(P \rho P)) = P_j \mathcal{E}_1(P \rho P) P_j - P_j \mathcal{E}_2(P \rho P) P_j
\]

\[
= \sum_{l \neq k} P_j F_l P \rho P F_k \dagger P_j - \sum_{l \neq k} P_j F_l P \rho P F_k \dagger P_j
\]

\[
= \sum_{l \neq k} P_j F_l P \rho P F_k \dagger P_k P_j - \sum_{l \neq k} P_j F_l P \rho P F_k \dagger P_l P_j.
\]

Note that \(\mathcal{E}_2(P \rho P) \neq 0\) because \(F_l P \rho P F_k \dagger \neq 0\) since \(F_l\) is a unitary matrix and thus preserves rank. Using the orthogonality of the \(P_i\) projector, we can choose \(j\) to be one of the \(l\) in the map so that \(P_j P_l = P_l\) for one \(l\) and the other \(P_i\) terms vanish. Thus, the probability of the outcome is a negative value. Therefore, enforcing \(\mathbf{27}\) results in a density matrix which has a negative eigenvalue and the code space is not in the domain of \(\mathcal{E}(\rho)\). \(\square\)

Note, however, that if \(\mathcal{E}_2(P \rho P) = 0\), we can still satisfy the NCP error correcting conditions and our code space is in the domain of \(\mathcal{E}\). This is stated more formally in the following theorem.

**Theorem 2.** Consider an evolution of the form \(\mathbf{19}\). If the quantum error correcting code conditions are satisfied, and \(\mathcal{E}_2(P \rho P) = 0\), i.e., the negative part of the map is zero on the code space, then the evolution can be reversed and the resulting density operator will be positive.

**Proof.** The proof follows from Theorem \(\mathbf{1}\) and the QEC code conditions for a CP map. Starting from \(\mathbf{23}\), we have (unnormalized)

\[
\mathcal{M}_j(\rho') = P_j \mathcal{E}_1(P \rho P) P_j - P_j \mathcal{E}_2(P \rho P) P_j
\]

\[
= P_j \mathcal{E}_1(P \rho P) P_j
\]

\[
= \sum_k |d_{kk}| P_j P_k U_k \rho U_k \dagger P_k P_j.
\]

From the orthogonality of \(P_i\), we have \(P_j P_k = P_k\) for one \(k\) value and the other \(P_i\) terms vanish. The correction is finished by conjugating with \(U_k \dagger\) because \(U_k \dagger P_k U_k = P_k\). This recovery process is given by the recovery map

\[
\mathcal{R}(\rho) = \sum_j U_j \dagger P_j \rho P_j U_j.
\]

Since \(\rho' = \mathcal{E}(P \rho P) = \mathcal{E}_1(P \rho P)\), \(\rho'\) is clearly positive. \(\square\)
IV. EXAMPLES

It is argued in Ref. [11] that, given an NCP map Φ, a corresponding CP map ˜Φ can be defined by taking the absolute value of the coefficients in the operator-sum decomposition. Then, using this CP map, a code space and recovery map is determined, which also works for the original NCP map. According to our Theorem 1, this can lead to a non-positive outcome, which we show with an example. In the Ref. 11, it states

Corollary 1. Consider a Hermitian noise map Φ_H(ρ) = \sum_{i=1}^{N} c_i K_i ρ K_i^† and associate to it a CP map Φ_{CP}(ρ) = \sum_{i=1}^{N} |c_i| K_i ρ K_i^†. Then any QEC code C and corresponding CP recovery map R_{CP} for Φ_{CP} are also a QEC code and CP recovery map for Φ_H.

The following gives an example of when their Corollary 1. produces a non-positive outcome which is covered by the corollary to Theorem 1.

Consider the three qubit bit flip map, used as the example in [11],

Φ_{IBF}(ρ) = c_0 ρ + c_1 \sum_{n=1}^{3} X_n ρ X_n,  \quad (31)

where X_n is the σ_x Pauli matrix acting on the n^{th} qubit, and c_0 and c_1 are real, have opposite sign, and c_0 + 3c_1 = 1. The corresponding CP map is Φ_{CP}(ρ) = |c_0|ρ + |c_1| \sum_{n=1}^{3} X_n ρ X_n, the code space is C = span\{|000⟩, |111⟩\}, and the projector onto the code space is P = |000⟩⟨000| + |111⟩⟨111|. Then,

R_{CP}[Φ_{IBF}(PρP)] ∝ PρP,  \quad (32)

where R_{CP} (given below) is the CP recovery map for Φ_{CP}. However, it turns out that the code space is not in the domain of the error map 31 and thus performing 32 leads to negative probabilities. This is what our Corollary 1 predicts.

Let

ρ = a |000⟩⟨000| + (1 - a) |111⟩⟨111| + c^∗ |111⟩⟨000| + c |000⟩⟨111|,  \quad (33)

be a valid arbitrary density matrix in the code space. Applying the error map 31 onto 33, we get

Φ_{IBF}(ρ) = c_0 [a |000⟩⟨000| + (1 - a) |111⟩⟨111| + c^∗ |111⟩⟨000| + c |000⟩⟨111]| + c_1 [a |100⟩⟨100| + (1 - a) |011⟩⟨011| + c^∗ |011⟩⟨100| + c |100⟩⟨011| + c_1 [a |010⟩⟨010| + (1 - a) |101⟩⟨101| + c^∗ |101⟩⟨010| + c |010⟩⟨101| + c_1 [a |001⟩⟨001| + (1 - a) |110⟩⟨110| + c^∗ |110⟩⟨001| + c |001⟩⟨110| + c^∗ |010⟩⟨010| + c |100⟩⟨100| + c^∗ |100⟩⟨000| + c |000⟩⟨100| = ρ ′,  \quad (34)

If we measure with projectors in the computational basis, the probabilities are given by tr(Pρ'). Then, the |000⟩⟨000|, |111⟩⟨111|, |100⟩⟨100|, and |011⟩⟨011| outcomes have corresponding probabilities tr(|000⟩⟨000|ρ') = c_0 a, tr(|111⟩⟨111|ρ') = c_0 (1 - a), tr(|100⟩⟨100|ρ') = c_1 a, and tr(|011⟩⟨011|ρ') = c_1 (1 - a). If ρ is a valid density matrix then it should be positive semidefinite and these values should be greater than or equal to zero. Here, regardless of which c_i is negative (as in the definition of Φ_{IBF}(ρ)), one of the resulting probabilities is negative. Thus, ρ' is not positive semidefinite and the code space C = span\{|000⟩, |111⟩\} is not in the domain of Φ_{IBF}(ρ).

Remark. We should emphasize here that we assume that the recovery map occurs after the error map so the output of the error map needs to be valid. However, it may be possible to implement the error and recovery maps together. In the latter situation, the code space would not need to be in the domain of the error map.

The recovery map for Φ_{CP}(ρ) is

R_{CP}(ρ) = PρP + \sum_{n=1}^{3} PX_nρX_nP  \quad (35)

If we apply this recovery to 34, we see that we get back to the initial state ρ. One may suppose that this works on average, but the processes of measurement, followed by a recovery, is nonphysical.

Remark. Shabani and Lidar [11] consider the Hermitian maps to be physical. Thus, the negativity of the outcome is not regarded. For example, later in Corollary 2, they consider a Hermitian recovery without regard to its positivity [11].

V. SUMMARY/DISCUSSION

In this paper, we address the reversibility of quantum operations for the evolution of a subsystem that does not correspond to a completely positive map. Some researchers suppose this is possible for a system that is initially correlated with its environment. The effects and reversibility of these more general error models were considered by both Aharonov and Ben-Or 10, and also Shabani and Lidar 11.

In general, we find that there are restrictions on the applicability of the standard quantum error correcting code conditions for evolutions that are not describable by a CP map if one is to expect a positive outcome for the operators. These restrictions are described in our theorem that shows that the diagonal CP error correcting conditions can fail to give a code space that has a positive output for these evolutions.

As a corollary, we also showed that when the pseudounitary degree of freedom diagonalizes the NCP error correcting condition and the operators in the diagonalized error map are unitaries, then the code space is not in the domain of the error map in the sense that it is not
positive. This implies that the applicability of the quantum error correcting conditions for linear maps given in [11] must be supplemented to achieve a positive outcome.

Finally, we presented a set of sufficient conditions for the reversibility of a more general error. This was followed by examples. In the near future, we will present other possibilities for reversing a more general quantum evolution.

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Appendix A: Properties of PseudoHermitian Matrices

Definition 1. A matrix is pseudounitary (PU) if
\[ U^\dagger = \eta U^{-1} \eta^{-1}, \tag{A1} \]
where \( \eta \) is a Hermitian matrix.

Definition 2. A matrix is pseudohermitian (PH) if
\[ H^\dagger = \eta H \eta^{-1}. \tag{A2} \]

Lemma 2. A PU matrix can be obtained from the exponentiation of a PH matrix.

Proof. Let \( U = \exp(-iHt) \) with \( H \) PH. Then
\[ U^\dagger = \exp(iH^\dagger t) = \exp(\eta H \eta^{-1} t) = \eta \exp(iHt) \eta^{-1} = \eta U^{-1} \eta^{-1}. \tag{A3} \]

Lemma 3. If a PU matrix is obtained from a matrix \( H \) via \( \exp(-iHt) \), then \( H \) is PH.

Proof. To see this, consider that
\[ U^{-1} U = I = \eta^{-1} U^\dagger \eta U, \]
so letting \( U = \exp(-itH) \),
\[ \frac{d}{dt} U^{-1} U \bigg|_{t=0} = \eta^{-1} iH^\dagger \eta - iH, \]
so \( H \) is PH.

As mentioned in the text, in this article, the \( \eta \) considered is of the form \( \eta = \text{diag}(1, 1, 1, \ldots, 1, -1, -1, \ldots, -1) \), with \( p \) ones and \( q \) minus ones. In this case, the form of the PH is
\[ \begin{pmatrix} A & B \\ -B^\dagger & C \end{pmatrix} \tag{A4} \]
where \( A \) is a \( p \times p \) Hermitian matrix, \( C \) is a \( q \times q \) Hermitian matrix and \( B \) is an arbitrary \( p \times q \) matrix. Note that this implies \( \eta H \) is Hermitian if \( H \) is PH and \( \eta M \) is PH, if \( M \) is Hermitian.

From the definition of a pseudounitary matrix, \( U \eta U^\dagger = \eta \) and for this unitary, the signature of the matrix \( \eta \) corresponds to the form of the unitary which is often denoted \( U(p,q) \) to emphasize this relation ship to \( \eta \) with the given “signature” \( p, q \).

Lemma 4. A pseudohermitian matrix \( H \) is diagonalizable by a matrix \( S \) via \( S^{-1} HS = H_d \), where \( H_d \) is diagonal and \( S \) can be chosen pseudounitary.

The following proof is adapted from Ref. [?] for the diagonalization of Hermitian matrices.

Proof. Let \( v_1 \) be an eigenvector of \( H \). (Every matrix has at least one eigenvector.) Let \( \lambda_1 \) be its corresponding eigenvalue. Then
\[ Hv_1 = \lambda_1 v_1. \]

Now we want to build a PU matrix that will diagonalize \( H \). Let \( v_1 \) be the first column of such a matrix and write
\[ U = \begin{pmatrix} v_1 & |Y \end{pmatrix}, \tag{A5} \]
where \( Y \) is an \( n \times (n - 1) \) matrix and \( v \) is an \( n \times 1 \) column vector. The matrix \( Y \) can be written as a set of \( n \times 1 \) column vectors \( v_i, i = 2, \ldots, n \). These vectors can be chosen orthogonal (under the \( \eta \) inner product) to \( v_1 \). (Or one could imagine using a Gram-Schmidt type process to make them orthogonal to \( v_1 \).) Thus,
\[ \langle \langle v_j, v_1 \rangle \rangle_{\eta} \equiv (v_j, \eta v_1) \equiv \sum_{k=1}^{n} (v_j^* k) \eta_k v_k = 0, \tag{A6} \]
where \( \eta_k \) are the diagonal elements of \( \eta \). (This could also be written using \( \eta_k = \eta \delta_{ik} \) and recall that \( \eta_k = \pm 1 \).) This implies that
\[ (Y^\dagger \eta v_1)_j = \sum_{k=1}^{n} Y^*_{kj} \eta_k v_k = 0. \tag{A7} \]
This is true for each \( j \), so \( Y^\dagger \eta v_1 = 0 \) as is \( \eta Y^\dagger v_1 = 0 \).

Now compute the following product
\[ U^\dagger \eta H U = \left( \begin{pmatrix} v_1^* \\ Y \end{pmatrix} \right) \eta H \left( \begin{pmatrix} v_1 & |Y \end{pmatrix} \right) \]
\[ = \left( \begin{pmatrix} v_1^* \eta v_1 & v_1^* \eta H Y \\ Y^\dagger \eta H v_1 & Y^\dagger \eta H Y \end{pmatrix} \right). \tag{A8} \]
This matrix has the following structure:

$$
\begin{pmatrix}
1 \times 1 & 1 \times (n-1) \\
(n-1) \times 1 & (n-1) \times (n-1)
\end{pmatrix}.
$$  \hfill (A9)

Now note that the upper left block and lower left block are

$$v_i^\dagger H v_1 = \lambda_1 v_i^\dagger \eta v_1 = \pm \lambda_1$$

$$Y^\dagger \eta H v_1 = \lambda_1 Y^\dagger \eta v_1 = 0.$$

\hfill (A10)

Now we have

$$U^\dagger \eta H U = \begin{pmatrix}
\pm \lambda_1 & 0 \\
0 & Y^\dagger \eta H Y
\end{pmatrix}.$$

Recall that $H$ is PH, so $H = \eta H^\dagger \eta$. This implies that

$$
(v_i^\dagger \eta H Y)^\dagger = Y^\dagger H^\dagger \eta v_1 = Y^\dagger \eta H^\dagger \eta v_1 = Y^\dagger \eta H v_1 = \lambda_1 Y^\dagger \eta v_1 = 0,
$$

where we have used the fact that $\eta^2 = 1$ and $\eta^\dagger = \eta$. At this point, we have that

$$U^\dagger \eta H U = \begin{pmatrix}
\pm \lambda_1 & 0 \\
0 & Y^\dagger \eta H Y
\end{pmatrix}.$$

\hfill (A11)

Now note that $N \equiv \eta U^\dagger \eta H U$ is PH since

$$\eta N^\dagger \eta = \eta (\eta U^\dagger \eta H U)^\dagger \eta = \eta (U^\dagger H^\dagger \eta U) \eta = U^\dagger \eta (\eta H^\dagger \eta) U = U^\dagger \eta H U = N,$$

where we have again used $\eta^2 = 1$ and $\eta^\dagger = \eta$. Also, note that $N$ clearly has the same form as $U^\dagger \eta H U$.

To see that $N$ has the same eigenvalues as $H$, we need only notice that $\eta U^\dagger \eta$ is $U^{-1}$ implies this since

$$U^{-1} H U w = \lambda w \Rightarrow H U w = \lambda U w.$$

\hfill (A15)

So that, letting $U w = v$, we see that for any $\lambda$

$$H v = \lambda v$$

\hfill (A16)

Thus, the eigenvalues of $N$ are the same as those of $H$. Notice that $\eta U^\dagger \eta = U^{-1}$ is exactly the PU condition that $U^\dagger \eta U = \eta$.

Given the form of the PH matrix, Eq. (A11), the matrix $\eta' Y^\dagger \eta H Y$ is also PH, where $\eta'$ is the same as $\eta$, albeit with one less diagonal entry. Thus, since $p$ and $q$ were arbitrary, this matrix can be treated in the exact same way as $H$. We can find an eigenvector and eigenvalue and reduce it in size by one leaving another PH matrix as a submatrix to be diagonalized. Continuing this allows the matrix to be diagonalized and the diagonalizing matrix is PU since $U^{-1} H U = \eta U^\dagger \eta H U$ is diagonalized.

\hfill \Box

Appendix B: PseudoUnitary Freedom in the Operator-Sum Representation

The Unitary degree of freedom for operators and for the operator-sum representation (OSR), is useful for a variety of reasons. The extension of the Unitary freedom for positive operators is extended to operators with negative eigenvalues. It is then shown that the freedom is also present in the OSR.

1. Unitary and PseudoUnitary Freedom for Operators

The unitary degree of freedom for operators is quite important since it shows that there are many different decompositions of a mixed state density operator [31]. This is discussed, for example, in textbooks [9, 32]. Refs. [33, 34] also provide interesting discussions and references. The non-uniqueness of a mixed state decomposition means that there are many different physical systems that could give rise to the same density operator (matrix).

The following is adapted from Nielsen and Chuang [9] with their Theorem stated below. Consider a density operator

$$\rho = \sum_i p_i |\psi_i\rangle \langle \psi_i| = \sum_i |\tilde{\psi}_i\rangle \langle \tilde{\psi}_i|,$$

where we define the unnormalized quantum state $|\tilde{\psi}_i\rangle \equiv \sqrt{p_i} |\psi_i\rangle$ and another decomposition of the same quantum state

$$\rho = \sum_j q_j |\phi_j\rangle \langle \phi_j| = \sum_j |\tilde{\phi}_j\rangle \langle \tilde{\phi}_j|,$$

where $|\tilde{\phi}_j\rangle \equiv \sqrt{q_j} |\phi_j\rangle$.

**Theorem 3.** (As stated in [2]. It is also proven there.) The sets $\{ |\tilde{\psi}_i\rangle \}$ and $\{ |\tilde{\phi}_j\rangle \}$ generate the same density matrix if and only if

$$|\tilde{\psi}_i\rangle = \sum_j u_{ij} |\tilde{\phi}_j\rangle,$$

where $(u_{ij})$ is a unitary matrix, and we add zero vectors to the smaller set so that the two sets have the same number of elements.

Now let us suppose that our operator can be expanded in a basis $\{ |v_i\rangle \}$ and a set of eigenvalues that are not necessarily positive, but are real, $\mu_i \in \mathbb{R}$,

$$\tau = \sum_i \mu_i |v_i\rangle \langle v_i|.$$

\hfill (B4)

Furthermore, suppose that there is another decomposition of $\tau$ in terms of a set of eigenvectors $\{ |w_j\rangle \}$ and eigenvalues $\nu_j$ so that we also have

$$\tau = \sum_j \nu_j |w_j\rangle \langle w_j|.$$

\hfill (B5)
As before, we define \( |\tilde{v}_i\rangle \equiv \sqrt{|\mu_i|} |v_i\rangle \) and \( |\tilde{w}_j\rangle \equiv \sqrt{|\mu_j|} |w_j\rangle \). We will also define \( \eta_i \equiv \text{sign}(v_i) \) and \( \nu_j \equiv \text{sign}(w_j) \) to be the sign (magnitude one) of the eigenvalues. Thus \( \eta_i = \pm 1 \); it is plus one for a positive eigenvalue and minus one for a negative eigenvalue.

**Theorem 4.** The sets \( \{ |\tilde{v}_i\rangle \} \) and \( \{ |\tilde{w}_j\rangle \} \) generate the same operator if and only if

\[
|\tilde{v}_i\rangle = \sum_j u_{ij} |\tilde{w}_j\rangle , \tag{B6}
\]

where \( (u_{ij}) \) is a pseudounitary matrix, and we add zero vectors to the smaller set so that the two sets have the same number of elements.

**Proof.** (\( \Leftarrow \)) Suppose \( |\tilde{v}_i\rangle = \sum_j u_{ij} |\tilde{w}_j\rangle \), where \( (u_{ij}) \) is a PU matrix. A number of zero vectors will be added to the smaller set to make them the same size. Therefore, the two sets of \( \eta_i \) and \( \nu_j \) can also be made the same size. Let us call them both \( \eta_i \). Then

\[
\tau = \sum_i \eta_i |\tilde{v}_i\rangle \langle \tilde{v}_i| = \sum_{ijk} \eta_{ij} \delta_{jk} \eta_k |\tilde{w}_j\rangle \langle \tilde{w}_j|,
\]

and since \( u_{ij} \) is PU, \( \sum_i u_{ij} \eta_i u_{ik}^* = \delta_{jk} \eta_k \). Therefore

\[
\tau = \sum_i \eta_i |\tilde{v}_i\rangle \langle \tilde{v}_i| = \sum_{ijk} \eta_k \delta_{jk} |\tilde{w}_j\rangle \langle \tilde{w}_j|,
\]

\[
= \sum_k \eta_k |\tilde{w}_k\rangle \langle \tilde{w}_k| . \tag{B8}
\]

(\( \Rightarrow \)) Now suppose

\[
\tau = \sum_i \mu_i |v_i\rangle \langle v_i| = \sum_j \nu_j |w_j\rangle \langle w_j|. \tag{B9}
\]

Let \( \tau = \sum_r \beta_r |\tilde{r}\rangle \langle \tilde{r}| \) be another decomposition of \( \tau \) with \( \{|\tilde{r}\rangle\} \) a complete set of unnormalized orthogonal states and \( \beta_r = \pm 1 \). The set of \( \{|\tilde{r}\rangle\} \) is complete, so we can append zeroes to the set \( \{|v_i\rangle\} \) and can take \( \beta_r = (\eta_i) \). Also since the set \( \{|\tilde{r}\rangle\} \) is complete, we can expand any \( |\tilde{v}_i\rangle \) as

\[
|\tilde{v}_i\rangle = M_{ir} |\tilde{r}\rangle . \tag{B10}
\]

Now, since these are both decompositions of \( \tau \), we have

\[
\tau = \sum_i \eta_i |\tilde{v}_i\rangle \langle \tilde{v}_i| = \sum_{irs} \eta_i M_{ir} M_{rs}^* |\tilde{r}\rangle \langle \tilde{s}| = \sum_r \eta_r |\tilde{r}\rangle \langle \tilde{r}| , \tag{B11}
\]

which is true if \( \eta_i M_{ir} M_{rs}^* = \eta_r \delta_{rs} \). This is just the condition for \( M \) to be pseudounitary.

Now, we could make the same argument for the decomposition in terms of \( |\tilde{w}_j\rangle \). Then, since these are each related by a PU and the composition of two PU matrices is a PU matrix, there exists a PU matrix that takes \( |\tilde{v}_i\rangle \) to \( |\tilde{w}_j\rangle \). \( \square \)

## 2. Unitary and PseudoUnitary Freedom in the OSR

The description of the dynamical map is not unique. It can be represented by the set of \( C_k \) corresponding to the eigenvector decomposition of the map \( B \), but there are many other representations. In this section we find an equivalence class of maps and provide an expression of such a freedom after reviewing the case for completely positive maps.

For completely positive maps, we reiterate that a theorem describing the freedom, examples, and uses can be found in Ref. [9].

### a. Unitary Freedom for Completely Positive Maps

Let us first quote Nielsen and Chuang [9]:

Suppose \( \{ E_1, ..., E_m \} \) and \( \{ F_1, ..., F_n \} \) are operation elements giving rise to quantum operations \( \mathcal{E} \) and \( \mathcal{F} \) respectively. By appending zero operators to the shorter list of operation elements we may ensure that \( m = n \). Then \( \mathcal{E} = \mathcal{F} \) if and only if there exist complex numbers \( u_{ij} \) such that \( E_i = \sum_j u_{ij} F_j \), and \( u_{ij} \) is an \( m \) by \( m \) unitary matrix.

Note that zero may be added to the map \( \mathcal{F} \) in such a way that it is not obtainable from the map \( \mathcal{E} \) by a unitary transformation. Let us consider the following example. Let \( \mathcal{E} \rightarrow \mathcal{E}' = \sum_i E_i E_i^\dagger + AA^\dagger - AA^\dagger \). Suppose \( A \) is linearly independent of all \( E_i \), then this map can not be obtained from the set with a unitary transformation. The map \( \mathcal{E}' \) differs from \( \mathcal{E} \) in some sense trivially and in practice it is very often easy to spot such an extension by zero.” However, the difference could be difficult to recognize and provides a technical point to note about the theorem.

When considering such cases, one may define an equivalence class of maps by identifying all maps which differ by such trivial extensions. Thus, maps which are in the same class are those which differ by the addition of operation elements which would cancel. The representative will always be the element of the class that has no such trivial extension. This will be termed a base map.

### Definition 3. For a given equivalence class of maps which differ by a trivial extension, the base map of the class is the representative of that class which has not been trivially extended.

Different base maps belong to different classes.

### b. Pseudo-unitary freedom for Hermiticity-preserving maps

Now let us consider a map \( \Phi(\rho) = \sum_j \eta_j C_j \rho C_j^\dagger \) and introduce a set of operators \( D_j \) corresponding to another
base map \( \Phi'(\rho) = \sum_j \eta_j D_j \rho D_j^\dagger \). As stated above, we may take \( \eta_j = \pm 1 \). We can choose the number of operators to be the same by appending zero operators to the shorter list. This enables the number of \(-1\) and \(+1\) to be chosen to be the same for each of the maps. Furthermore, we will order the set of \( \eta_j \) such that the first \( p \) are \(+1\) and the next \( q \) are \(-1\).

The freedom in the operator-sum representation is described by the group \( U(p,q) \). This group is often called a pseudo-unitary group due to its relation to the unitary group and it is a metric-preserving group with the signature of the metric determined by the integers \( p, q \). See for example (32, pages 45, 197), (36, page 392), (37, page 12), or (38, page 444).

Let \( \eta \) be an \( N \times N \) diagonal matrix with the first \( p \) entries \(+1\), the next \( q \) entries \(-1\), and \( N = p + q \). Then for all \( U \in U(p,q) \),

\[
U^\dagger \eta U = \eta U^{-1}.
\]

We may express the matrix \( \eta \) as a diagonal matrix with the matrix elements being \( \eta_k \), \( \eta_k = +1 \), for \( k = 1, \ldots, p \) and \( \eta_k = -1 \), for \( k = p + 1, \ldots, p + q = N \). Alternatively, we may express the matrix \( \eta \) using elements \( (\eta)_{kl} = \eta_k \delta_{kl} \). This is a diagonal matrix with the first \( p \) entries along the diagonal are \(+1\) and the next \( q \) are \(-1\).

Let the elements of the matrix \( U \) be given by \( u_{ij} \) and those of \( U^\dagger \) be \( u^*_{ij} \). Then the Eq. (B12) can be written as \( U^\dagger \eta U = \eta \), or since \( \eta^2 = I \), \( U \eta U^\dagger = \eta \). In components, this can be written as

\[
\sum_{jk} u_{ij} \eta_j \delta_{jk} u^*_{ik} = \eta_i \delta_{il}.
\]

Having established this property for elements of the group \( U(p,q) \), the following theorem may now be stated and proved. (Originally, a version of the following proof was presented in Ref. [30].)

**Theorem 5. Pseudo-unitary freedom:** Suppose \( \{C_1, C_2, \ldots, C_n\} \) and \( \{D_1, D_2, \ldots, D_m\} \), are operation elements giving rise to base quantum operations (maps) \( \Phi \) and \( \Phi' \) respectively. Explicitly,

\[
\Phi = \sum_i \gamma_i C_i C_i^\dagger, \quad \Phi' = \sum_j \mu_j D_j D_j^\dagger,
\]

where each \( \gamma_i \) and each \( \mu_j \) is \( \pm 1 \) and ordered as above, with all \(+1\) eigenvalues first. Furthermore, we can always take \( \gamma_i = \mu_i \) with zero-valued \( C_i \) or zero-valued \( D_j \) appended to the shorter list for the \(+1\) (\(-1\) eigenvalue. Then \( \Phi = \Phi' \) if and only if

\[
D_j = \sum_i u_{ji} C_i,
\]

where the numbers \( u_{ji} \) form a \( p + q \) by \( p + q \) matrix in \( U(p,q) \).

**Proof:** We first consider whether the condition is necessary and use the notation \( C_i = |i\rangle \), \( D_i = |j\rangle \). Suppose that

\[
\Phi = \Phi'.
\]

(Or, if one would like to display the argument explicitly, \( \Phi(\rho) = \Phi'(\rho) \).) For a general map \( \Phi \), there exists a corresponding \( B \) matrix such that \( \Phi = B \) (i.e. \( \Phi(\rho) = B \rho \)). \( B \) has an eigenvector decomposition \( B = \sum_k \lambda_k^\dagger |k'\rangle \langle k'| \) where the set of \( |k'\rangle \) are linearly independent since they are orthogonal. This follows from the fact that the eigenvectors can be chosen orthogonal. Now \( |k\rangle = \sqrt{|\lambda_k|} |k'| \). These vectors are clearly also orthogonal and thus linearly independent if the \( |k'\rangle \) are. Then \( B \) can be re-expressed as \( B = \sum_k \eta_k |k\rangle \langle k| \) with the first \( p \) eigenvalues \( \eta_k = +1 \), \( k = 1, \ldots, p \) and the next \( q \) eigenvalues \( \eta_k = -1 \), \( k = p + 1, \ldots, p + q \). This gives

\[
B = \sum_k \eta_k |k\rangle \langle k| = \sum_{k=1}^p |k\rangle \langle k| - \sum_{k=p+1}^{p+q} |k\rangle \langle k|,
\]

which is an eigenvector decomposition of the map \( \Phi \). Now, let us consider another decomposition of \( B \) corresponding to the set of \( C_i \), \( B = \sum \gamma_i |i\rangle \langle i| \). Each \( |i\rangle \) can be written as a linear combination of the \( |k\rangle \), \( |i\rangle = \sum_k w_{ik} |k\rangle \). (See for example Ref. [9], page 104.) Given \( \Phi = B \)

\[
\sum_k \eta_k |k\rangle \langle k| = \sum_{kl} \left( \sum_i \gamma_{i} w_{ik} w_{il}^* \right) |k\rangle \langle l|.
\]

Since the \( |k\rangle \) are linearly independent, it is clear that this can only happen if

\[
\sum_i \gamma_{i} w_{ik} w_{il}^* = \delta_{kl} \eta_k.
\]

We may always take \( \eta_k = \gamma_{i} \) by appending the shorter list of vectors \( \{|i\rangle\} \) or \( \{|k\rangle\} \) with zero vectors. This will ensure the matrices \( \gamma \) with elements \( \delta_{ij} \gamma_{i} \) and \( \eta \) with elements \( \delta_{ij} \eta_{k} \) are equal. Furthermore, \( w \) can then be taken to be square with \( |i\rangle = \sum_k w_{ik} |k\rangle \). The condition, Eq. (B19), can then be written as

\[
w^\dagger \eta w = \eta,
\]

which is the condition for the matrix \( w \) to be in \( U(p,q) \). Now, we can use the same argument, with \( B = \Phi' \) and \( v_{jk} \) such that \( |j\rangle = \sum_k v_{jk} |k\rangle \), to show

\[
v^\dagger \eta v = \eta.
\]

Since each of these two are related to the same expression for \( B \) using elements of \( U(p,q) \) which is a group, then the linear transformation which takes the \( C_i \) to the \( D_j \), \( u = vw^{-1} \) is in \( U(p,q) \).
Next, we consider whether \( u \in U(p, q) \) will imply that \( \Phi = \Phi' \), i.e., if the condition is sufficient. This is straightforward algebra. Then, Eq. (135)

\[
\Phi'(\rho) = \sum_j \mu_j D_j \rho D_j^\dagger = \sum_{lkj} \mu_j \mu_{kj} u_{jk}^* C_l \rho C_l^\dagger
\]

\[
= \sum_{lkj} \left( \sum_j \mu_j \mu_{kj}^* u_{jk} \right) C_l \rho C_l^\dagger
\]

\[
= \sum_l \gamma_l \delta_{lk} C_l \rho C_l^\dagger = \Phi(\rho),
\]

which shows that the two sets of operators \( C_j \) and \( D_j \) related by a pseudo-unitary matrix \( u \) will yield the same map. \( \square \)

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