Artificial Intelligence Based A* Optimization Routing in Mobile Ad Hoc Networks

Gondi Yasoda Devi, Gurnala Venkateswara Rao

ABSTRACT: A Mobile Ad Hoc Network (MANET) is a hotchpotch of nodes with mobility feature, the established network utilization is dynamically outlined based on temporary architecture. In MANETs, the challenging and vital role is played by the routing protocols performance factors under different condition and environments. The routing protocols are liable to handle many nodes with limited resources. There exit many routing protocols in MANETs, one of the main key note that has to be considered in designing a routing protocol is to observe that the designed routing protocol is having an proportionate effect on network performance. The existence of obstacles may lead to many geographical routing problems like excess consumption of power and congestion of data. The aim of this paper is to take the assistance of A* algorithm that finds the walk-able path avoiding the concave obstacle in the path relaying on the gaming-theory model[29]. This algorithm decreases the delays in packet transmission and in turn increases the success rate of transmission. We take into consideration path length, penalty for node availability as probability of forwarding criteria and processes effective packet transmission. The simulated results analyse the performance of our protocol over other conventional algorithms based on congestion cost, path length, node availability penalty, delay, packet loss, throughput.

KEYWORDS: MANETS, A*, Penalty of Node Availability, Path Length, Heuristic.

I. INTRODUCTION

Wireless communication is tomost communication channel of this era form top to bottom in any part of the world. Several challenges has to be taken into consideration while enhancing the technology. One of the most dynamic technology we are very fond off, in communication, without which we feel like our hand are cut is the Mobile adhoc networks (MANETs), where a researcher has to pay a key role in designing the optimistic routing protocols. In wireless communication MANETs are self regulating automated networks. In MANETs the nodes are by nature having the feature of mobility they do not relay on infrastructure and nodes communicate and interact with each other across the wireless the interfaces [9]. MANETs are easily applicable when there is a requirement for temporary communication system without infrastructure for instance events of large wild and region stricken by earthquakes [10]. The nodes communicate with each other within the given cluster or nodes within another cluster heads directly. Hence, in networking each node accomplish client and route functionalities simultaneously [11].

The design of the routing algorithms must integrate several factors owing to the MANETs features, e.g. node’s trust, save energy, network topology dynamic changes. Considering the trust of node, the interaction of multi-hop nodes is relied on the node’s consistency of the route. Therefore, the existing node’s trustworthiness is recognized significantly in the routing protocols. In MANETs, the capacity of the mobile node's battery is limited, which affects the survival of the network because the links are not connected while the battery is worn out. On considering the routing protocol, the energy of the mobile nodes is necessary for ensuring the network connectivity and enhances the life span of the network.

Factors that play a critical role in routing protocol

A. Optimal and Effective Routing: Which is one of the primary requirement of the defined protocol. We have to consider this as the primary goal of the designed protocol has to unearth and transfer the packets from the source to the destination. Packet Delivery Ratio is one of the parameter for efficient routing, end-to-end delay’s average and optimal rout’s percentage are also considered to make the routing protocol efficient and optimal.

B. Congestion Avoidance: It Strongly evaluated that congestion is linked with packet drops, if the congestion is more the packet drop rate will be more which decreases the efficiency of routing protocol.

C. Energy Consumption: Due to the increased smart devices and mobile nodes like PDA’s, laptops, cellular phones and other portable devices are having strong demand towards energy consumption. To the analysis made by experts individual nodes energy consumption should be minimized.

D. Load Balancing: The nodes that located in the communication of optimal routes may suffer from overloads which might lead to network congestion. Fail distribution of load has to be attained between the individual nodes in optimal routes.

E. Reachability: The probability of finding at least a single path between the selected source and destination nodes is considered as reachability.

During routing, the flexibility of the nodes made it complex to identify the network topology utilized by nodes. Moreover, the volatile environment, the untrustworthy mobile nodes, and the wireless medium might cause different defects in MANETs. There are several suggestions on different strategies and protocols, as the “Internet Engineering Task Force and academic and industrial undertakings” are taking additional standardization efforts [24].

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The conventional routing algorithm is categorized into three groups, such as proactive, hybrid, reactive routing on the basis of time triggered route discovery [13] [25]. Every node in the defined network will have global view and routing table is updated in certain time intervals through every node communicate its routing table with neighbours. In proactive routing, every node communicates its route table occasionally for its neighbours to the extent that each node is having a global view of the network, for example PSR [14], and FSR [15]. Nodes create routes only on request in reactive routing such as AODV [23], and CAODV [16]. In general, hybrid routing separates the network into various regions, and then proactive routing is selected in the region and reactive routing between the regions, like ZRP [17], and HOPNET [13]. Excessive emphasis on routing trust leads to improve the control overhead and decrease the routing efficiency. At the same time overemphasis on routing efficiency refuses the additional features that might react to routing efficiency. Few existing methodologies are implemented to enhance the efficiency and routing trust i.e., cryptography and data hiding, location-aware and their combinations, trust value, and detection.

The major contribution of the paper is to frame an optimized A* algorithm. A* is one among the best optimization search algorithm in artificial intelligence. It makes use of heuristic values for path-finding. A* searching technique can be considered as one of the algorithm that follows smart technique which make it better than any conventional algorithms. The algorithm has been used to discover shortest path in some of the game designing and maps that are web-based very efficiently. The existence of obstructing obstacles in the network area challenges the design of routing protocols in both static and mobile ad hoc networks. In A* Algorithm at each and every time it selects the node based on the value of ‘T’ which equal to the added value of h and g. The node selection is done at every step by finding the node with lowest ‘T’. The ‘g’ and ‘h’ values are defined as:

\[ g = \text{cost incurred in forming a path from the source node to destination in the selected square of the grid.} \]

\[ h = \text{cost incurred in forming a path from selected square grid, i.e., the current node to the destination node, we consider heuristic as a smart guess.} \]

The organization of the paper is designed according to the details mentioned below:

- Section II reviews the features and challenges of conventional MANETs routing algorithms.
- Section III speaks about routing strategy in MANET: network and system model with proposed routing.
- Section IV in this section Obstacle aware routing pattern in MANET using A* algorithm is given.
- Section V illustrates simulation setup of the system
- Section VI Performance evaluation of routing protocols
- Section VII concludes this paper
One of the primary challenges in framing a MANET is enabling each device to incessantly sustain the data transmission without any traffic.

Since MANETs are a type of wireless ad hoc networks, it usually posses a routable networking environment on top of a Link Layer ad hoc network. Moreover, the routing strategy often suffers from some obstacles, which need to be focussed for the development of improved routing desires. Hence, it is highly necessary to find a novel structure the wireless network. Assume \( M \) number of mobile nodes is uniformly distributed in the network with area \( MD \times ND \).

**C. Proposed Flow Graph for Obstacle Aware Routing in MANET**

In the past years, many researchers had dedicated their precious time for improving the efficiency of MANETs routing algorithms. The main intend of this paper is to develop obstacle-aware routing protocol in MANET using optimized A* algorithm. The flow diagram of the proposed MANET routing is shown in Fig. 3.

![Diagrammatic representation of proposed MANET Routing](image)

**Fig. 3: Diagrammatic representation of proposed MANET Routing**

A* algorithm, which is a well-performing path selection algorithm by concerning the obstacles. Even though this algorithm is optimally efficient to solve complex routing problems, its performance has to be still improved by taking into account few more constraints other than node availability and distance. Hence, in order to further improve the performance of shortest route path selection by A* algorithm, an improved meta-heuristic algorithm termed as CE-CSO is deployed for optimally placing the obstacles in different locations for finding the optimal shortest path. Here, the improved shortest path selection is done in such a way that the total cost concerning the sum of the distance between the path and the congestion of each node in the network is minimized.

**III. OBSTACLE AWARE ROUTING PATTERN IN MANET USING A* ALGORITHM**

The obstacle aware shortest path is determined by the A* algorithm. In general, A* algorithm [29] is used to find the shortest paths from the source to destination by removing obstacles. Moreover, it is designed with a classical grid-based approach. For defining the process of A* to find the shortest path, the diagrammatic representation is given in Fig. 4.

![Shortest path representation](image)

**Fig. 4: Shortest path representation**

In Fig. 4, the D the target node, obstacles are indicated by the black node, and the source node is denoted by the S, and the other blocks are for the normal nodes. In the above diagram, A* algorithm chooses each step in an efficient manner i.e., the algorithm travels from (Ro1, Co1) to (Ro3, Co3) not to (Ro2, Co3). In addition, the algorithm moves from (Ro3, Co3) to (Ro4, Co2) but not to (Ro4, Co3). In the mean time, the squares have distance attribute separately consisting of \( F(b) \), \( G(b) \), and \( H(b) \) values. The cost \( F(b) \) is the sum of \( G(b) \) and \( H(b) \), which is mathematically represented in Eq. (1). The value of \( G(b) \) is the movement cost from source node to the destination node \( b \) and the value of \( H(b) \) is the movement cost from the current node \( b \) to the destination node(b).

\[
F(b) = G(b) + H(b)
\]  
(1)

In order to reach the destination node from the source node in a shortest way, A* algorithm is utilized. By using A* algorithm, the minimum \( F \) value need to be found out. This procedure needs to be repeated until the shortest path is found with minimum cost.
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It is happened by maintaining the paths starting at the source node and moving towards the path end at a time until the termination criterions met.

In Eq. (1), \( b \) is the next node in the path. When A*selects the path to extend from source node to the target node, the A* search is terminated. For performing repeated selection of minimum cost nodes to extend, the applications of A* makes use of priority queue. Moreover, the priority queue is called open set or fringe.

Here, at each step, the node consisting of low \( F(b) \) value will be eradicated from the queue. The values of \( F(b) \) and \( G(b) \) of its neighbors are updated accordingly and those are added to the queue. The A* continues its search until the destination node has minimum \( F(b) \) value than any other node in the queue, and the corresponding value will be the cost of shortest path.

In an admissible heuristic, the value of \( H(b) \) is zero at the destination node. In order to define the exact series of steps, the algorithm is easily revised, thus every node in the path has the information of predecessor node. This algorithm will run up to the destination node by pointing to the predecessor and so on, until few node’s predecessor is the source node.

**A* Algorithm :**

Two list are defined as CLOSED list and OPEN list

1. Open list is initialized
2. Closed list is initialized
   
   In the open list take the starting node(by marking ‘f’ as zero)
3. Loop begins by checking for open list is empty or not

   4. Node’s least ‘f’ value from open is selected, let it be ‘q’

   5. Pop ‘q’

   6. ‘n’ q predecessor are generated

   7. For every predecessor
      i) if the goal is predecessor, search is stopped
         predecessor.g = q.g + (distance from q to predecessor)
         predecessor.h = distance from predecessor to goal
         predecessor.f = predecessor.g + predecessor.h

      ii) the predecessor is skipped, if the value of ‘f’ is less than predecessor at same position from OPEN list

      iii) the predecessor is skipped, if the node of CLOSED list at same position as of predecessor has a lower f value than predecessor, else to OPEN list add that node

   end (for)

8. Into closed list q is pushed

end (loop)

**Euclidean Distance :**

- Heuristic value is the distance form the current node to destination node is considered as distance formula

\[
H = \sqrt{(\text{current_node.x} - \text{destination.x})^2 + (\text{current_node} - \text{destination.y})^2}
\]

- To find the path the signals are allowed to walk in any route, we can make use of heuristic values

The determination of obstacle ware shortest path is done by the A* algorithm, in which the node availability and distance between the path is considered. Here, the obstacle is recognized by the CE-OLSR [35]. The dynamic and automatic detection of obstacles that present in the network is accomplished by CE-OLSR with no extra signalling overhead. It is an advanced version of OLSR protocol, which depends on two concepts like “network cartography, and stability routing”. In order to build an accurate network topology, CE-OLSR utilizes network cartography as a substitute to link states. For avoiding the usage of weak links during the routing procedure, CE-OLSR depends on stability routing.

**IV. SIMULATION SETUP OF THE SYSTEM**

The simulation setup involves network basic settings the simulation is performed in Matlab simulator.

The below table gives the considered parameter values for the proposed A* algorithm

| Table I : Routing Parameters |
|-----------------------------|
| Parameters | Value |
| Number of nodes | Varies |
| Mac Type | 802.11e |
| Protocol | AODV, A* |
| Communication Protocol | UDP |
| Application | CBR |
| Delay | 1ms |
| Simulation Time | Varies |

**V. PERFORMANCE EVALUATION OF ROUTING PROTOCOLS**

This design focuses on the specified metrics performance and are measured quantitatively. The accuracy analyzed over the designed algorithm, performance metrics plays a main role. The metrics that are considered are as follows,

**A. Delay:** The average time spend for the transmission of a packet from source to destination point. The Total_Delay_of_Packets at Destination_Node is \( d \), and the Total_Number_of_Packets_Received at the destination is \( pkt_d \), n’s Traffic iThe results of the proposed algorithm

**B. Path Length:** Path length is considered as the distance that is measured from the source node to the destination node.

**C. Congestion Cost:** Congestion cost is considered as the cost that is charged due to the problem of congestion

**D. Penalty for Node Availability:** The main criteria of laying the path is the availability of nodes. if more time is taken in finding the availability of the node, the over-all outcome may be effected. this delay in identifying the neighboring nodes is termed as penalty of node availability. Path finding time is effected by the node penalty.
E. Packet Loss: when the data packet are transmitted through the network path in selected route when any packet doesn’t reach the destination point i.e the packet is dropped in the middle before the reach of destination then the loss of packet occurs. Packet loss is also proportional to the congestion in the network. Loss of Packet loss is calculated as the percentage of lost packets of sent packets.

F. Throughput: The count of bits transmitted per unit of time from source point to the destination point.

VI. SIMULATION RESULTS

About MATLAB simulator

For technical computations and attain high-performance the language MATLAB can be used. We can find integration of visualization, as an easy computation environment, visualization and mathematical notations can be might lead to the solutions that are familiar. The other uses can discussed as the, Development of algorithm, Computation and Math, Simulation, prototyping, modelling the system, And many more

The below result are obtained by implementation of the A* algorithm in the MATLAB

While considering the network with obstacle the how the signal propagation is being taken place and based on which the route formed is depicted in below figures 5 and 6 for AODV and A* algorithms respectively.

The below tables depict the performance of two routing protocol i.e., the conventional routing protocol AODV and proposed obstacle aware optimal artificial intelligence routing protocol A*, where the network of light and dense are being considered.

The following simulation results elucidates the comparison between conventional AODV and A* algorithm. When AODV algorithm is applied for packet routing the congestion cost and packet loss are minimalistic when compared to A* algorithm. But the other parameters such as path length and penalty of node availability, delay show efficiency when implanted with A* rather than with conventional AODV algorithm. The overall cost calculated by considering all the given parameters in the table of packet routing through A* is optimal when compared to AODV

TABLE II. PERFORMANCE ANALYSIS OF PROPOSED AND CONVENTIONAL OBSTACLE-AWARE MANET ROUTING FOR 78, 108, 128, 200 NODES

| Measures For 78 Nodes | AODV  | A*   |
|-----------------------|-------|------|
| Path Length           | 176.29| 67.52|
| Congestion Cost       | 15.882| 25.638|
| Penalty for Node Availability | 160 | 0 |
| Delay                 | 0.71  | 0.35 |
| Packet Loss           | 0.1   | 0.34 |
| Throughput            | 126.76| 94.286|
| Total Cost            | 352.99| 93.859|

| Measures For 108 Nodes | AODV  | A*   |
|-----------------------|-------|------|
| Path Length           | 194.29| 130.03|
| Congestion Cost       | 2.1272| 40.416|
| Penalty for Node Availability | 220 | 0 |
| Delay                 | 0.71  | 0.68 |
| Packet Loss           | 0.08  | 0.69 |
| Throughput            | 129.58| 22.794|
| Total Cost            | 417.22| 171.86|
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| Measures For 128 Nodes | AODV | A* |
|------------------------|------|----|
| Path Length            | 217.36 | 186.39 |
| Congestion Cost        | 9.269  | 44.007 |
| Penalty for Node Availability | 220   | 0    |
| Delay                  | 0.71   | 0.61 |
| Packet Loss            | 0.07   | 0.62 |
| Throughput             | 130.99 | 31.148 |
| Total Cost             | 447.41 | 231.66 |

| Measures For 200 Nodes | AODV | A* |
|------------------------|------|----|
| Path Length            | 247.45 | 217.22 |
| Congestion Cost        | 16.943 | 37.1 |
| Penalty for Node Availability | 170  | 0 |
| Delay                  | 0.65   | 0.74 |
| Packet Loss            | 0.06   | 0.75 |
| Throughput             | 144.62 | 16.892 |
| Total Cost             | 435.11 | 255.87 |

![Graphs showing comparisons between AODV and A* for different measures](image-url)
VII. CONCLUSION

The A* signalling, can be used for automatic presumption of the occurrences of static obstacles positioned in the network. This work identifies the obstacles and lays the shortest path optimally. Simulation results reveals that defined metrics improves the accuracy and coverage ratio. The attained results shows that the proposed algorithm improves the effectiveness by optimally recognising the obstacle area and minimizing the path laying cost and penalty of node availability. The proposed work provides a sufficient accuracy ratio to economically full fill our aim to autonomously avoid broken links due to obstacle obstruction. The proposed obstacle detection algorithm is integrated with CE-OLSR

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