A GENERALIZATION OF THE PROBLEM OF MARIUSZ MESZKA

ANITA PASOTTI AND MARCO ANTONIO PELLEGRINI

Abstract. Mariusz Meszka has conjectured that given a prime \( p = 2n + 1 \) and a list \( L \) containing \( n \) positive integers not exceeding \( n \) there exists a near 1-factor in \( K_p \) whose list of edge-lengths is \( L \). In this paper we propose a generalization of this problem to the case in which \( p \) is an odd integer not necessarily prime. In particular, we give a necessary condition for the existence of such a near 1-factor for any odd integer \( p \). We show that this condition is also sufficient for any list \( L \) whose underlying set \( S \) has size 1, 2, or \( n \). Then we prove that the conjecture is true if \( S = \{1, 2, t\} \) for any positive integer \( t \) not coprime with the order \( p \) of the complete graph. Also, we give partial results when \( t \) and \( p \) are coprime. Finally, we present a complete solution for \( t \leq 11 \).

1. Introduction

Throughout this paper \( K_v \) will denote the complete graph on \( \{0, 1, \ldots, v-1\} \) for any positive integer \( v \). For the basic terminology on graphs we refer to [17]. Following [12], we define the length \( \ell(x, y) \) of an edge \([x, y]\) of \( K_v \) as

\[
\ell(x, y) = \min(|x - y|, v - |x - y|).
\]

If \( \Gamma \) is any subgraph of \( K_v \), then the list of edge-lengths of \( \Gamma \) is the multiset \( \ell(\Gamma) \) of the lengths (taken with their respective multiplicities) of all the edges of \( \Gamma \). The set of the edges of \( \Gamma \) will be denoted by \( E(\Gamma) \). Also, by \( \delta(\Gamma) \) we will mean the multiset

\[
\delta(\Gamma) = \{|x - y| : [x, y] \in E(\Gamma)\}.
\]

Clearly, if all the elements of \( \delta(\Gamma) \) do not exceed \( \frac{v}{2} \) it results \( \delta(\Gamma) = \ell(\Gamma) \). For our convenience, if a list \( L \) consists of \( a_1 \) 1's, \( a_2 \) 2's, \ldots, \( a_t \) \( t \)'s, we will write \( L = \{1^{a_1}, 2^{a_2}, \ldots, t^{a_t}\} \), whose underlying set is the set \( \{1, 2, \ldots, t\} \).

The following conjecture [5, 18] is due to Marco Buratti (2007, communication to Alex Rosa).

Conjecture (Buratti). For any prime \( p = 2n + 1 \) and any list \( L \) of \( 2n \) positive integers not exceeding \( n \), there exists a Hamiltonian path \( H \) of \( K_p \) with \( \ell(H) = L \).

In [12] Peter Horak and Alex Rosa generalized Buratti’s conjecture. Such a generalization has been restated in a easier form in [13] as follows.

Conjecture (Horak and Rosa). Let \( L \) be a list of \( v - 1 \) positive integers not exceeding \( \lfloor \frac{v}{2} \rfloor \). Then there exists a Hamiltonian path \( H \) of \( K_v \) such that \( \ell(H) = L \) if, and only if, the following condition holds:

\[
\text{for any divisor } d \text{ of } v, \text{ the number of multiples of } d \text{ appearing in } L \text{ does not exceed } v - d.
\]
Let

Proof. condition (2) we have to prove that $|v| \leq 23$. With the aid of a computer, we have verified its validity for all odd integers not exceeding 23. We recall that a

Recently, Mariusz Meszka formulated a very similar conjecture concerning near 1-factors. We recall that a near 1-factor of $K_{2n+1}$ is a subgraph of $K_{2n+1}$ consisting in $n$ disjoint edges and one isolated vertex, while a 1-factor of $K_{2n}$ is a subgraph of $K_{2n}$ consisting in $n$ disjoint edges.

Conjecture (Meszka). For any prime $p = 2n + 1$ and any list $L$ of $n$ positive integers not exceeding $n$, there exists a near 1-factor $F$ of $K_p$ with $\ell(F) = L$.

To best of our knowledge, all the results on the problem proposed by Meszka are contained in [15], where Rosa proved that the conjecture is true when the elements of the list are $\{1, 2, 3\}$ or $\{1, 2, 3, 4\}$ or $\{1, 2, 3, 4, 5\}$. Also he proved that the conjecture is true when the list $L$ contains a sufficient number of 1’s, in detail, if $L = \{1^{a_1}, 2^{a_2}, \ldots, n^{a_n}\}$ with $a_1 \geq \frac{n}{2}$. Moreover, with the aid of a computer, we have verified the validity of Meszka’s conjecture for all primes $p \leq 23$.

Working on this conjecture it is easy to see that the assumption $p$ prime is not necessary, as it happens for Buratti’s one. Indeed we found another condition which is necessary and which led us to propose our conjecture.

Conjecture. Let $v = 2n + 1$ be an odd integer and $L$ be a list of $n$ positive integers not exceeding $n$. Then there exists a near 1-factor $F$ of $K_v$ such that $\ell(F) = L$ if, and only if, the following condition holds:

$$\text{for any divisor } d \text{ of } v, \text{ the number of multiples of } d \text{ appearing in } L \text{ does not exceed } \frac{v-d}{2}. \tag{2}$$

With the acronym MPP, which stands for Meszka-Pasotti-Pellegrini, we will denote this more general conjecture. In particular MPP($L$) will denote the conjecture for a given list $L$. Clearly, if $v$ is a prime then our conjecture reduces to Meszka’s one. With the aid of a computer, we have verified its validity for all odd integers $v \leq 23$. We point out that in the statement, the actual conjecture is the sufficiency. In fact we can prove that condition (2) is necessary.

Proposition 1. The list $L$ of edge-lengths of any near 1-factor of $K_v$ satisfies condition (2).

Proof. Let $F$ be a near 1-factor of $K_v$ with $\ell(F) = L$ and let $d$ be a divisor of $v$. Denote by $D$ the sublist of $L$ containing all the multiples of $d$ appearing in $L$, hence we have to prove that $|D| \leq \frac{v-d}{2}$. Note that if the length of an edge is a multiple of $d$, the vertices of such an edge have to be in the same residue class modulo $d$ and that $F$ has exactly $\frac{v}{d}$ vertices in each residue class modulo $d$. Also, $\frac{v}{d}$ is an odd number and this implies that with the elements of each residue class modulo $d$ we can construct at most $\frac{v-1}{2}$ edges whose length is in $D$. Since there are exactly $d$ residue classes modulo $d$ we have at most $\frac{v-d}{2}$ edges whose length is in $D$, namely $|D| \leq \frac{v-d}{2}$. \hfill \Box

Before giving the main result of this paper we would like to show some connections between MPP-problem and graph decompositions, as done for BHR-problem in [13]. For a general background on graph decompositions see [3]. Reasoning in the same way as we have done in [13] for the BHR-problem, one can easily obtain that MPP($L$) can be reformulated in the following way.
Conjecture. A Cayley multigraph $Cay[\mathbb{Z}_n : A]$ admits a cyclic decomposition into near 1-factors if and only if $A = L \cup -L$ with $L$ satisfying condition (2).

For reader convenience we recall the definition of a Cayley multigraph. A list $\Omega$ of elements of an additive group $G$ is said to be symmetric if $0 \notin \Omega$ and the multiplicities of $g$ and $-g$ in $\Omega$ coincide for any $g \in G$. If $\Omega$ does not have repeated elements then one can consider the Cayley graph on $G$ with connection set $\Omega$, denoted $Cay[G : \Omega]$, whose vertex-set is $G$ and in which $[x, y]$ is an edge if and only if $x - y \in \Omega$. Cayley graphs have a great importance in combinatorics and they are precisely the graphs admitting an automorphism group acting sharply transitively on the vertex-set (see, e.g., [11]). If, more generally, the symmetric list $\Omega$ has repeated elements one can consider the Cayley multigraph on $G$ with connection multiset $\Omega$, also denoted $Cay[G : \Omega]$ and with vertex-set $G$, where the multiplicity of an edge $[x, y]$ is the multiplicity of $x - y$ in $\Omega$ (see, e.g., [6]).

In the next sections of this paper we prove some results concerning the MPP-problem that we can summarize in the following theorem.

Theorem 2. Let $L$ be a list of $n$ positive integers not exceeding $n$ and let $v = 2n + 1$. Then, $MPP(L)$ holds whenever the underlying set $S$ of $L$ satisfies one of the following conditions:

1. $|S| = 1, 2$ or $n$;
2. $S = \{1, 2, t\}$, where $t \geq 3$ is not coprime with $v$.

Furthermore, $MPP(L)$ holds if $L = \{1^a, 2^b, t^c\}$ with $a + b \geq \left\lceil \frac{t-1}{2}\right\rceil$.

The proof of Theorem 2 will follow from Propositions 11, 14, 16 and 18.

If $S = \{1, 2, t\}$ and $\gcd(v, t) = 1$, we also present some partial results for the cases not covered by Theorem 2. In particular, we give a complete solution for $t \leq 11$.

2. Linear realizations and their relationship with Skolem sequences

A cyclic realization of a list $L$ of $n$ positive integers not exceeding $n$ is a near 1-factor $F$ of $K_{2n+1}$ such that $\ell(F) = L$. For example, given $L = \{1^2, 4^3, 6\}$ the near 1-factor $F = \{(0, 4), (1, 2), (5, 12), (6, 10), (7, 11), (8, 9)\} \cup \{3\}$ of $K_{13}$ is a cyclic realization of $L$. Clearly if a list $L$ admits a cyclic realization we can say that $MPP(L)$ is true.

A linear realization of a list $L$ with $n$ positive integers not exceeding $2n$ is a near 1-factor $F$ of $K_{2n+1}$ such that $\delta(F) = L$. It is quite immediate that if the elements of $L$ do not exceed $n$ a linear realization of $L$ is nothing but a cyclic realization, namely a near 1-factor $F$ of $K_{2n+1}$ with $\delta(F) = \ell(F) = L$. Following [15], we will say that a linear realization of a list $L$ is almost perfect or perfect if the isolated vertex of the near 1-factor is $2n - 1$ or $2n$, respectively.

It is important to underline that there is a strong relationship between (almost) perfect linear realizations and Skolem sequences. We point out that Skolem sequences and their generalizations (see, e.g., [1]) have revealed to be very useful in the construction of several kinds of combinatorial designs (see, for example, the survey [10] and the references therein as well as [2, 4, 19]). In order to present this connection we recall the basic definitions, see [16]. A Skolem sequence of order $n$ is a sequence $S = (s_1, s_2, \ldots, s_{2n})$ of $2n$ integers satisfying the following conditions:

1. for every $k \in \{1, 2, \ldots, n\}$ there exist exactly two elements $s_i, s_j \in S$ such that $s_i = s_j = k$;
It is worth to observe that a Skolem sequence can also be written as a collection of ordered pairs \( \{(a_i, b_i) : 1 \leq i \leq n, b_i - a_i = i\} \) with \( \bigcup_{i=1}^{n}\{a_i, b_i\} = \{0, 1, \ldots, 2n - 1\} \). For instance, the Skolem sequence \( S = (1, 1, 3, 4, 5, 3, 2, 4, 2, 5) \) of order 5 can be seen as the set \( \{(0, 1), (6, 8), (2, 5), (3, 7), (4, 9)\} \). Hence we can conclude that a Skolem sequence is equivalent to a perfect linear realization of a set. For example, \( S' = (1, 1, 3, 4, 5, 3, 2, 4, 2, 5, 0) \), obtained from \( S \) adding 0 at the end, is a perfect linear realization of the set \( \{1, 2, 3, 4, 5\} \). Also, a hooked Skolem sequence of order \( n \) is a sequence \( HS = (s_0, s_1, \ldots, s_{2n}) \) of \( 2n + 1 \) integers satisfying above conditions (1) and (2) and such that \( s_{2n-1} = 0 \). So we have that a hooked Skolem sequence is an almost perfect linear realization of a set.

Hence, we can say that MPP-problem can be view as a generalization of Skolem sequences and we propose the following more general definition. Let \( L = \{1^{a_1}, 2^{a_2}, \ldots, n^{a_n}\} \) be a list on the set \( \{1, 2, \ldots, n\} \) and let \( k \) be an element of the set \( \{1, 2, \ldots, 2n + 1\} \). We call \( k \)-extended Skolem sequence of \( L \) any sequence \( S = (s_0, s_1, \ldots, s_{2n}) \) for which it is possible to partition \( \{1, 2, \ldots, 2n + 1\} \) \( \setminus \{k\} \) into a set \( T \) of \( n \) ordered pairs \( (x, y) \) with \( x < y \) such that the set \( T_i := \{(x, y) \in T \mid s_x = s_y = y - x = i\} \) has size \( a_i \) for \( 1 \leq i \leq n \). It is clear that a linear realization of \( L \) is a \( k \)-extended Skolem sequence of \( L \) for a suitable \( k \). Also, it is perfect or almost perfect when \( k = 2n \) or \( 2n - 1 \), respectively.

In these cases, in view of the classical definitions given above of Skolem sequences, one may speak of an ordinary or hooked Skolem sequence of \( L \), respectively. For example, given \( L_1 = \{1, 3, 6^2\} \) the near 1-factor \( F_1 = \{[0, 6], [1, 7], [3, 4], [5, 8]\} \cup \{2\} \) of \( K_9 \) is a linear realization of \( L_1 \). The corresponding 2-extended Skolem sequence is \( S_1 = (6_1, 6_2, 0, 1, 1, 3, 6_1, 6_2, 3) \), where we use \( 6_1 \) and \( 6_2 \) to distinguish the same length 6 belonging to distinct pairs. Take now the near 1-factor \( F_2 = \{[0, 3], [1, 4], [2, 6]\} \cup \{5\} \) of \( K_7 \); it is an almost perfect realization of \( L_2 = \{3^2, 4\} \) and it is easy to see that the corresponding hooked Skolem sequence is \( S_2 = (3_1, 3_2, 4, 3_1, 3_2, 0, 4) \).

Clearly a list \( L \) can admit both a not perfect linear realization and a perfect linear realization. For instance \( F_3 = \{[0, 6], [1, 7], [2, 5], [3, 4]\} \cup \{8\} \) is a perfect linear realization of the above list \( L_1 = \{1, 3, 6^2\} \), which corresponds to the Skolem sequence \( S_3 = (6_1, 6_2, 3, 1, 1, 3, 6_1, 6_2, 0) \). It is easy to see that if there exists a perfect linear realization of a list \( L \), then there exists a 1-factor \( F \) of \( K_{2|L|} \) such that \( \delta(F) = L \).

For convenience, in the following, by \( rL \), \( apL \) and \( pL \) we will denote a linear realization, an almost perfect linear realization and a perfect linear realization of \( L \), respectively.

Given two lists \( L_1 \) and \( L_2 \) it is possible to obtain a linear realization of the list \( L_1 \cup L_2 \) composing the linear realizations of \( L_1 \) and \( L_2 \). We have to point out that it is not always possible to compose two linear realizations, since the existence of the composition depends on the properties of the two realizations as shown in the following lemma, see [15]. When the composition exists it will be denoted by “+”.

Lemma 3. Given two lists \( L_1 \) and \( L_2 \) we have:

1. \( pL_1 + pL_2 = p(L_1 \cup L_2) \);
2. \( pL_1 + apL_2 = ap(L_1 \cup L_2) \);
3. \( apL_1 + apL_2 = p(L_1 \cup L_2) \);
4. \( pL_1 + rL_2 = r(L_1 \cup L_2) \).
Proof. (1) (2) (3) These items have been proved by Rosa, see [15].
(4) Let $|L_1| = w$ and $|L_2| = z$. Let $S_1 = (\ell_1, \ell_2, \ldots, \ell_{2w}, 0)$ be the sequence corresponding to a perfect linear realization of $L_1$ and let $S_2 = (\tilde{\ell}_1, \tilde{\ell}_2, \ldots, \tilde{\ell}_i, 0, \tilde{\ell}_{i+1}, \ldots, \tilde{\ell}_{2z})$ be the sequence corresponding to a linear realization of $L_2$. Then $S = (\ell_1, \ell_2, \ldots, \ell_{2w}, \tilde{\ell}_1, \tilde{\ell}_2, \ldots, \tilde{\ell}_i, 0, \tilde{\ell}_{i+1}, \ldots, \tilde{\ell}_{2z})$ is the sequence corresponding to a linear realization of $L_1 \cup L_2$. \hfill \Box

Example 4. Consider the previous linear realizations of the lists $L_1 = \{1, 3, 6^2\}$ and $L_2 = \{3^2, 4\}$. Using Lemma 3 one gets a perfect realization of $\{1^2, 3^2, 6^4\}$, an almost perfect realization of $\{1, 3^4, 4, 6^2\}$, a perfect realization of $\{3^4, 4^2\}$ and a linear realization of $\{1^2, 3^2, 6^4\}$. Indeed, $p\{1^2, 3^2, 6^4\} = pL_1 + pL_1$ can be obtained using twice the sequence $S_1$ to get the sequence $(61, 62, 31, 11, 11, 31, 61, 62, 63, 64, 32, 12, 12, 32, 63, 64, 0)$. Using the sequences $S_1$ and $S_2$ one gets the sequence $(61, 62, 31, 11, 11, 31, 61, 62, 32, 33, 4, 32, 33, 0, 4)$, corresponding to $ap\{1, 3^3, 4, 6^2\} = pL_1 + apL_2$. Now, using twice the sequence $S_2$, one gets the sequence $(3, 3, 4, 3, 1, 3, 3, 4, 2, 4, 1, 3, 3, 3, 3, 4, 3, 4, 0)$ corresponding to $p\{3^4, 4^2\} = apL_2 + apL_2$. Finally, using the sequences $S_3$ and $S_1$, one gets the sequence $(61, 62, 31, 11, 11, 31, 61, 62, 32, 33, 0, 12, 12, 32, 63, 63, 32)$, corresponding to $r\{1^2, 3^2, 6^4\} = pL_1 + rL_1$.

Given a list $L$ and a positive integer $q$, by $q \cdot pL$ we will mean the perfect linear realization $\underbrace{pL + pL + \ldots + pL}_{q \text{ times}}$.

In view of the above lemma we will look for linear realizations possibly (almost) perfect.

Corollary 5. There exists a perfect linear realization of $L = \{(2i + 1)^{a_i} \mid i = 0, \ldots, n, \forall n \in \mathbb{N}, a_i \geq a_j \text{ if } i < j\}$.

Proof. Given an odd integer $2k + 1$, it is immediate that $S_k = (2k + 1, 2k - 1, 2k - 3, \ldots, 9, 7, 5, 3, 1, 1, 3, 5, 7, 9, \ldots, 2k - 3, 2k - 1, 2k + 1, 0)$ corresponds to a perfect linear realization of the list $L_k = \{1, 3, 5, \ldots, 2k - 1, 2k + 1\}$. Hence, by Lemma 3 a linear realization of $L$ is $pL = a_n \cdot pL_n + \sum_{i=0}^{n-1}(a_i - a_i+1) \cdot pL_i$. \hfill \Box

Lemma 6 (A. Rosa, [15]). The list $\{x^x\}$ admits a perfect linear realization for each $x \geq 1$.

Lemma 7. The list $L = \{x^{x-1}\}$ admits a linear realization for each $x \geq 1$. If $x = 2$, such a realization is almost perfect.

Proof. The near 1-factor $F = \{i, i + x \mid 0 \leq i \leq x - 2\} \cup \{x - 1\}$ of $K_{2x-1}$ is such that $\delta(F) = L$. If $x = 2$ the statement immediately follows. \hfill \Box

Corollary 8. If there exists a linear realization of a list $L = \{\ell_1^{a_1}, \ell_2^{a_2}, \ldots, \ell_n^{a_n}\}$ then there exists a linear realization of the list $L' = \{\ell_1^{k_1+\ell_1}, \ell_2^{k_2+\ell_2}, \ldots, \ell_n^{k_n+\ell_n}\}$ for any $k_1, k_2, \ldots, k_n \in \mathbb{N}$. Also, if the realization of $L$ is (almost) perfect, the realization of $L'$ has the same property.

Proof. Let $rL$ be a linear realization of $L$. By Lemma 6 there exists $p\{\ell_i^{k_i}\}$ for any $\ell_i \in L$. A linear realization of $L'$ is $rL' = \sum_{i=1}^{n} k_i \cdot p\{\ell_i^{k_i}\} + rL$, see Lemma 3 (4). If $rL$ is almost perfect or perfect the thesis follows from Lemma 3 (2) and (1), respectively. \hfill \Box
Hence \( \text{MPP}(\{x, y\}) \) and \( pL_1 \) corresponding to the sequence \( S_3 = \{6, 1, 3, 2, 1, 1, 3, 6, 2, 0\} \). Then, for instance, \( p(1, 2^4, 3^4, 6^2) = 2 \cdot p(2^2) + p(3^4) + pL_1 \) exists and can be obtained taking the sequence \( (2, 2, 2, 1, 2, 3, 2, 4, 3, 1, 3, 3, 6, 1, 6, 2, 3, 4, 1, 1, 3, 6, 1, 6, 2, 0) \).

3. First cases

Let \( L \) be a list of \( n \) positive integers not exceeding \( n \) and let \( S \) be its underlying set. We start investigating \( \text{MPP}(L) \) in the following cases: \( |S| = 1, 2 \) and \( n \).

**Remark 10.** Let \( F \) be a near 1-factor of \( K_{2n+1} \) with \( \ell(F) = \{x_1^{a_1}, x_2^{a_2}, \ldots, x_n^{a_n}\} \). Let \( y \) be an integer coprime with \( 2n + 1 \) and let \( r_i \) be the remainder of the division of \( yx_i \) by \( 2n + 1 \). Define \( \ell_i = r_i \) if \( 1 \leq r_i \leq n \) and \( \ell_i = 2n + 1 - r_i \) if \( n < r_i \leq 2n \).

Then, multiplying each vertex of \( F \) by \( y \), it is possible to obtain a near 1-factor \( F' \) of \( K_{2n+1} \) such that \( \ell(F') = \{\ell_1^{a_1}, \ell_2^{a_2}, \ldots, \ell_n^{a_n}\} \).

If \( |S| = 1 \), all the edges have the same length. If this length is 1 it is immediate to see that \( F = \{[2i, 2i + 1] \mid i = 0, \ldots, n - 1\} \cup \{2n\} \) is a near 1-factor \( F \) of \( K_{2n+1} \) such that \( \ell(F) = \{1^n\} \). Let now \( S = \{x\} \) with \( 2 \leq x \leq n \). By Proposition 1 we have to consider only the case \( \gcd(x, 2n + 1) = 1 \). By Remark 10, multiplying the vertices of \( F \) by \( x \), we obtain a near 1-factor \( F' \) of \( K_{2n+1} \) such that \( \ell(F') = \{x^n\} \). Hence we can conclude that \( \text{MPP}(\{x^n\}) \) is true for any positive integers \( n \) and any positive integer \( x \) not exceeding \( n \).

If \( |S| = n \), namely if \( L = S = \{1, 2, 3, \ldots, n\} \), a near 1-factor of \( K_{2n+1} \) such that \( \ell(F) = L \) is, up to translations, a starter of \( Z_{2n+1} \). In fact, a starter in the odd order abelian group \( G \) (written additively), where \( |G| = 2n + 1 \), is a set of unordered pairs \( R = \{\{r_i, t_i\} \mid 1 \leq i \leq n\} \) that satisfies:

1. \( \{r_i \mid 1 \leq i \leq n\} \cup \{t_i \mid 1 \leq i \leq n\} = G \setminus \{0\} \)
2. \( \{\pm(r_i - t_i) \mid 1 \leq i \leq n\} = G \setminus \{0\} \)

Hence \( \text{MPP}(\{1, 2, 3, \ldots, n\}) \) is always true for any positive integer \( n \), in fact it is sufficient to take \( F = \{[i, 2n - 1 - i] \mid i = 0, \ldots, n - 1\} \cup \{2n\} \), namely, the so-called patterned starter of \( Z_{2n+1} \), see [8].

So, we proved the following.

**Proposition 11.** Let \( L \) be a list of \( n \) elements not exceeding \( n \) with underlying set \( S \). Then, \( \text{MPP}(L) \) holds if either \( |S| = 1 \) or \( |S| = n \).

Now, we consider the case \( |S| = 2 \), i.e. \( L = \{x^n, y^b\} \), where \( 0 < x < y \leq a + b = n \). Observe that by Proposition 1 we have to consider only the case \( \gcd(x, y, 2n + 1) = 1 \).

We start considering the case \( x = 1 \).

**Lemma 12.** There exists a linear realization of any list \( L = \{1^n, y^b\} \) whenever \( a \geq \lfloor \frac{y}{2} \rfloor \).

**Proof.** Set \( v = 2(a + b) + 1 \) and write \( b = qy + r \), where \( 0 \leq r < y \). Firstly, we consider the following sequences. If \( y - r \) is even, take

\[
S_1 = (y_1, y_2, \ldots, y_r, 1, 1, 1, 2, 1_2, \ldots, 1_{x-r}, 1_{x-r}, y_1, y_2, \ldots, y_r, 0)
\]

and if \( y - r \) is odd, take

\[
S_2 = (y_1, y_2, \ldots, y_r, 1, 1, 1, 2, 1_2, \ldots, 1_{x-r-1}, 1_{x-r-1}, 0, y_1, y_2, \ldots, y_r).
\]
Case 2. gcd(v, y)

hence, there exists a near 1-factor

Let a < b and if the proof in four cases. But firstly it is important to observe that if j and x ≥ ⌊L−1⌋, we have found a perfect realization of L. If b is odd, our realization is almost perfect.

It follows from previous corollary that, given a list L = \{1^a, 2^b\} there exists a near 1-factor F of K_v, where v = 2(a + b) + 1, such that δ(F) = L whose isolated vertex is v − 2 if b is odd, v − 1 if b is even. In the following, we will denote by F + g, g ∈ N, the graph obtained from F by replacing each vertex x of F with x + g. Clearly, F + g is a near 1-factor of the complete graph whose vertex set is \{g, i + g, \ldots, 2(a + b) + g\} such that δ(F + g) = L. This remark will be very useful in the next section.

Now we are ready to deal with the general case of two distinct edge lengths.

**Proposition 14.** MPP(\{x^a, y^b\}) holds for any a, b, x, y > 0.

Proof. Let L = \{x^a, y^b\} with x, y ≤ a + b and let v = 2(a + b) + 1. We have to split the proof in four cases. But firstly it is important to observe that if a ≥ b then a ≥ ⌊\frac{a+1}{2}\], since y ≤ a + b.

Case 1. gcd(x, v) = gcd(y, v) = 1. If a ≥ b we multiply the elements of L by x−1 and if a < b, we multiply by y−1. In both cases, we obtain a list L′ with underlying set S′ = \{1, z\} and z ≤ a + b. So, L′ satisfies the assumption of Lemma 12 and hence, there exists a near 1-factor F′ such that ℓ(F′) = L′. By Remark 10, there exists a near 1-factor F of K_v such that ℓ(F) = L.

Case 2. gcd(v, y) = d_2 > 1 and a ≤ \frac{(d_2−1)v}{2d_2}. Clearly it has to be gcd(x, d_2) = 1. Also by Proposition 1, a ≥ \frac{d_2−1}{2}.

We write the vertices of K_v, namely the integers 0, 1, \ldots, v − 1, within a matrix M of d_2 rows and \frac{v}{d_2} columns, where the element m_{i,j} is (i−1)x + (j−1)y modulo v, for i = 1, \ldots, d_2 and j = 1, \ldots, \frac{v}{d_2}. It is not hard to see that all the integers m_{i,j} are distinct. In fact if, by way of contradiction, m_{i,j} = m_{h,k}, then (i−1)x + (j−1)y ≡ (h−1)x + (k−1)y (mod v). In particular, this means that ix ≡ hx (mod d_2), namely d_2 | (i − h)x. Since gcd(x, d_2) = 1, it follows that d_2 | (i − h) and so i = h, since i, h ≤ d_2. Hence, (j−1)y ≡ (k−1)y (mod v). Then, v | (j−k)y and so \frac{x}{d_2} | (j−k). We obtain that \frac{x}{d_2} | (j−k), i.e. j = k, since j, k ≤ \frac{v}{d_2}.

Now, we construct the edges for the required near 1-factor using exactly once all but one the elements of the matrix M as vertices. Consider the following \frac{d_2−1}{2} edges of length x: \{m_{2i+1,1}, m_{2i+2,1}\} for i = 0, \ldots, \frac{d_2−3}{2}.

If a − \frac{d_2−1}{2} is even, let a − \frac{d_2−1}{2} = (\frac{d_2−1}{2})q + r with 0 ≤ r < \frac{d_2−1}{2} − 1, note that r is even. We take the following edges of length x: \{m_{2i+1,j}, m_{2i+2,j}\} for i = 0, \ldots, q − 1 and \{m_{2q+1,j}, m_{2q+2,j}\} for j = 2, \ldots, \frac{x}{d_2} (if q > 0) and \{m_{2q+1,j}, m_{2q+2,j}\} for j = 2, \ldots, r + 1. It is easy to see that, since r is even, there are an even number of elements which are not used to construct the above edges, in each row of M except for the last one. Hence we construct the b edges of length y connecting pairs of adjacent elements in each
rows. Since $a \leq \frac{(d_2-1)v}{2d_2}$ we have a sufficient number of edges of length $y$. In such a way we obtain a near 1-factor of $K_v$ whose isolated vertex is $d_{2,1}$.

If $a - \frac{d_2-1}{2}$ is odd, we construct the following edge of length $x$. Observing that $m_{d_{2,1}} + x = (d_2 - 1)x + x = d_2x$ is an element of the first row, say $m_{1,k}$, we can take the edge $[m_{d_{2,1}}, m_{1,k}]$. Hence we have to construct other $a - \frac{d_2-1}{2}$ edges of length $x$. Since $a - \frac{d_2+1}{2}$ is even, let $a - \frac{d_2+1}{2} = (\frac{d_2}{d_2} - 1)q + r$ with $0 \leq r < \frac{d_2}{2} - 1$, note that $r$ is even. We take the following edges of length $x$: \([m_{2i,j}, m_{2i+1,j}]\) for $i = 1, \ldots, q$ and $j = 2, \ldots, \frac{d_2}{2}$ (if $q > 0$) and \([m_{2q+2,j}, m_{2q+3,j}]\) for $j = 2, \ldots, r + 1$. It is easy to see that, since $r$ is even, there are an even number of elements, which are not used to construct the above edges, in each row of $M$ except for the first one. Hence, we can construct the edges of length $y$ as in the previous case.

Case 3. gcd$(v, y) = d_2 > 1$, $a > \frac{(d_2-1)v}{2d_2}$, gcd$(x, v) = 1$. By the assumption on $a$, it follows that $a > b$. In fact we have $a > \frac{(d_2-1)v}{2d_2} = \frac{d_2-1}{d_2} \geq \frac{2}{3} \frac{v}{3} = \frac{v}{3} = \frac{2(a + b) + 1}{3}$ and so $a > 2b + 1 > b$. Hence we can proceed analogously to Case 1, applying Remark 10 and Lemma 12.

Case 4. gcd$(v, y) = d_2 > 1$, $a > \frac{(d_2-1)v}{2d_2}$, gcd$(x, v) = d_1 > 1$. Note that in this case $b \leq \frac{(d_2-1)v}{2d_1}$. In fact, if, by way of contradiction, we have $b > \frac{(d_2-1)v}{2d_1}$ it results

$$\frac{v - 1}{2} = a + b > \frac{(d_1 - 1)v}{2d_1} + \frac{(d_2 - 1)v}{2d_2} = \frac{v}{2} \left( \frac{d_1 - 1}{d_1} + \frac{d_2 - 1}{d_2} \right) \geq \frac{v}{2} \left( \frac{2}{3} + \frac{2}{3} \right) = \frac{2v}{3} > \frac{v - 1}{2},$$

which clearly is a contradiction. Hence we can apply the same process of Case 2, interchanging $x$ with $y$. \hfill $\square$

**Example 15.** Take now $L = \{6^9, 10^{13}\}$. Here $v = 45$, $d_1 = \text{gcd}(6, 45) = 3$, $d_2 = \text{gcd}(10, 45) = 5$. Note that we are in Case 2 of Proposition 14, and that $a - \frac{d_2-1}{2} = 9 - 2 = 7$ is odd. Firstly, we construct the matrix

$$M = \begin{pmatrix}
0 & 10 & 20 & 30 & 40 & 5 & 15 & 25 & 35 \\
6 & 16 & 26 & 36 & 1 & 11 & 21 & 31 & 41 \\
12 & 22 & 32 & 42 & 7 & 17 & 27 & 37 & 2 \\
18 & 28 & 38 & 43 & 3 & 13 & 23 & 33 & 8 \\
24 & 34 & 44 & 9 & 19 & 29 & 39 & 4 & 14 \\
\end{pmatrix}.$$

Reasoning as in the proof of previous proposition we obtain the following edges of lengths 6: \{0, 6\}, \{12, 18\}, \{24, 30\}, \{16, 22\}, \{26, 32\}, \{36, 42\}, \{1, 7\}, \{11, 17\}, \{21, 27\}. The elements used to construct these edges are highlighted in bold in the matrix $M$. Now, we can construct the following edges of length 10: \{10, 20\}, \{40, 5\}, \{15, 25\}, \{31, 41\}, \{37, 2\}, \{28, 38\}, \{3, 13\}, \{23, 33\}, \{43, 8\}, \{34, 44\}, \{9, 19\}, \{29, 39\}, \{4, 14\}. So, the isolated vertex is 35.

4. **Near 1-factors with edge lengths $1, 2, t$**

In this section we investigate MPP($L$) where $L = \{1^a, 2^b, t^c\}$ for any integer $t > 2$. Clearly, in view of the results of the previous section, we can assume $a, b, c \geq 1$.

**Proposition 16.** Let $L = \{1^a, 2^b, t^c\}$ with $a + b \geq \lceil \frac{t+c}{2} \rceil$. Then there exists a linear realization of $L$. 
Proof. Let \( c = tq + r \) with \( 0 \leq r \leq t - 1 \). If \( r = 0 \) the statement follows from Lemmas 3 and 6 and Corollary 13. So we can assume \( r \geq 1 \). Since \( a + b \geq \lceil \frac{t}{2r} \rceil \) then \( a + b \geq \lfloor \frac{r}{t} \rfloor \). We start constructing a linear realization of \( L' = \{ t^a, 2^b, t^r \} \) and we have to distinguish four cases according to the congruence class of \( t - r \) modulo 4.

Case 1. \( t - r \equiv 0 \) (mod 4).

If \( b \geq \frac{t-r}{2} \), we consider the sequence \( S = (t, 2, 2, 2, 2, 2, 2, \ldots, 2, t, 2, 2, 2, 2, 2, 2, \ldots, 2, t, 2, 2, 2, 2, 2, 2, \ldots, 2, t, 2, 2, 2, 2, 2, 2, \ldots) \) which is a perfect realization of \( L'' = \{ 2^{\frac{t-r}{2}}, t^r \} \).

Next, consider \( rL' = pL'' + \{ 1^a, 2^b \} \).

Now, assume \( b < \frac{t-r}{2} \). If \( b \) is even, the sequence \( S = (t, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots) \) is a perfect linear realization of \( L'' = \{ 1^{t-r-2}, 2^b, t^r \} \). In this case \( pL' = pL'' + \{ 1^a, 2^b \} \).

If \( b \) is odd, the sequence \( S = (t, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots) \) is a perfect linear realization of \( L'' = \{ 1^{t-r-2}, 2^b, t^r \} \). Hence, \( \text{ap} L' = pL'' + \{ a, 2^b \} \).

Case 2. \( t - r \equiv 1 \) (mod 4).

Suppose \( b \) even. If \( b \geq \frac{t-r-2}{2} \), the sequence \( S = (t, 2, 2, 2, 2, 2, 2, \ldots, 2, t, 2, 2, 2, 2, 2, \ldots, 2, t, 2, 2, 2, 2, 2, 2, \ldots) \) is a linear realization of \( L'' = \{ 2^{t-r-2}, t^r \} \) and \( rL' = p\{ 1^a, 2^b \} + rL'' \). If \( b < \frac{t-r-2}{2} \), the sequence \( S = (t, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots) \) is a linear realization of \( L'' = \{ 1^{t-r-2}, 2^b, t^r \} \). In this case, \( rL' = p\{ 1^a, 2^b \} + rL'' \).

Suppose now \( b \) odd. If \( b \geq \frac{t-r-1}{2} \), the sequence \( S = (t, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots) \) is a linear realization of \( L'' = \{ 1^{t-r-2}, 2^b, t^r \} \). In this case, we obtain \( rL = p\{ 1^a, 2^b \} + rL'' \).

Case 3. \( t - r \equiv 2 \) (mod 4).

If \( b \geq \frac{t-r-2}{2} \), the sequence \( S = (t, 2, 2, 2, 2, 2, 2, \ldots, 2, t, 2, 2, 2, 2, 2, \ldots, 2, t, 2, 2, 2, 2, 2, 2, \ldots) \) is a perfect linear realization of \( L'' = \{ 1, 2^{t-r-2}, t^r \} \).

Hence, \( rL' = pL'' + \{ 2^{t-r-2}, t^r \} \).

Assume now \( b < \frac{t-r-2}{2} \). If \( b \) is even, we take the sequence \( S = (t, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots) \) that is a perfect linear realization of \( L'' = \{ 1^{t-r-2}, 2^b, t^r \} \). In this case, \( pL' = pL'' + \{ 1, 2^{t-r-2}, t^r \} \).

If \( b \) is odd, we take the sequence \( S = (t, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots, 2, 2, 2, 2, 2, 2, 2, \ldots) \) that is a perfect linear realization of \( L'' = \{ 1^{t-r-2}, 2^b, t^r \} \). So, \( \text{ap} L' = pL'' + \{ 1, 2^{t-r-2}, t^r \} \).

Case 4. \( t - r \equiv 3 \) (mod 4).

Suppose \( b \) even. If \( b \geq \frac{t-r-3}{2} \), the sequence \( S = (t, 2, 2, 2, 2, 2, 2, \ldots, 2, t, 2, 2, 2, 2, 2, \ldots, 2, t, 2, 2, 2, 2, 2, 2, \ldots) \) is a linear realization of \( L'' = \{ \ldots \} \).


\[\{1, 2^{\frac{b-1}{2}}, t^r\}\]. In this case, \(rL' = p\{1^{a-1}, 2^{b-\frac{t-1}{2}}\} + rL''\). If \(b < \frac{t-r-1}{2}\) the sequence \(S = (t_1, t_2, \ldots, t_r, 2, 1, 2, 1, 2, 1, 2, 2, 1, 2, 1, 2, 2, 1, \ldots, 2, 2, 1, 2, t_r, 0, t_1, t_2, \ldots, t_r)\) is a linear realization of \(L'' = \{1^{t-1-2b-1}, 2^b, t^r\}\). So, we have \(rL' = p\{1^{a-\frac{t-r-1}{2}}\} + rL''\).

Suppose now \(b \geq \frac{t-r-1}{2}\) the sequence \(S = (t_1, t_2, \ldots, t_r, 21, 2, 1, 2, 1, 2, 2, 1, 2, 1, 2, 2, 1, 2, 1, \ldots, 2, 2, 1, 2, t_r, 0, t_1, t_2, \ldots, t_r)\) is a linear realization of \(L'' = \{2^{\frac{t-1}{2}}\}, t^r\}\). Here, \(rL' = p\{1^{a}, 2^{b-\frac{t-1}{2}}\} + rL''\). If \(b < \frac{t-r-1}{2}\) the sequence \(S = (t_1, t_2, \ldots, t_r, 21, 2, 1, 2, 1, 2, 2, 1, 2, 1, 2, 2, 1, 2, 1, \ldots, 2, 2, 1, 2, t_r, 0, t_1, t_2, \ldots, t_r)\) is a linear realization of \(L'' = \{1^{\frac{t-1-2b-1}{2}}, 2^b, t^r\}\). Hence, \(rL' = p\{1^{a-\frac{t-r-1}{2}}\} + rL''\) is a linear realization of \(L' = \{1^a, 2^b, t^r\}\).

Now in order to obtain the thesis it is sufficient to note that \(rL = q \cdot p(t^t) + rL'\).

Example 17. Let \(L = \{1^4, 2^2, 12^{26}\}\). Since \(4 + 2 > \lfloor \frac{12-1}{2} \rfloor\) we are in the hypothesis of previous proposition. Note that in this case \(q = 2\) and \(r = 2\), hence \(t - r = 12 - 2 = 2\) (mod 4) so we are in Case 3. Also \(b = 2 < 4 = \frac{t-r-1}{2}\), so, since \(b\) is even, we take the following perfect linear realization of \(L' = \{1^3, 2^2, 12\}\): \(S = (12, 12, 2, 1, 2, 1, 2, 1, 1, 1, 2, 1, 1, 1, 2, 1, 1, 1, 2, 1, 1, 1, 2, 1, 1, 1, 2)\). Now a perfect linear realization of \(L\) is given by \(2 \cdot p(12^{12}) + p(1^{1}) + pL'\).

Proposition 18. MPP(L) holds for any \(L = \{1^a, 2^b, t^r\}\), with \(\gcd(t, v) = d > 1\), where \(v = 2(a + b + c) + 1\).

Proof. Clearly, \(t \leq \frac{a-1}{2}\) and by Proposition 1 we can assume \(a + b \geq \frac{d-1}{2}\). Also, in view of Proposition 16 we can assume \(a + b < \lfloor \frac{r+1}{2} \rfloor\) which implies \(c > \frac{r+1}{2} - \lfloor \frac{r+1}{2} \rfloor\), hence we have also \(c \geq \frac{a+b}{2}\). In fact, if, by way of contradiction, \(\frac{a+b}{2} - \lfloor \frac{r+1}{2} \rfloor < \frac{a+b}{2}\), we have \(\frac{a+b}{2} < \frac{r+1}{2}\), which cannot be \(d \geq 3\).

Now, let \(M = (m_{i,j})\) be the \(d \times \frac{v}{d}\) matrix whose elements are so defined, for \(i = 1, \ldots, d\) and \(j = 1, \ldots, \frac{v}{d}\):

\[m_{i,j} = \begin{cases} 
(i - 1) + (j - 1)t \pmod{v} & \text{if } i \equiv 0, 1 \pmod{4} \\
(i + (j - 1)t) \pmod{v} & \text{if } i \equiv 2 \pmod{4} \\
(i - 2) + (j - 1)t \pmod{v} & \text{if } i \equiv 3 \pmod{4}. 
\end{cases}\]

It is easy to see that the elements of \(M\) are the vertices of \(K_v\). Also, it is not hard to check that \(M\) has at least 5 columns. In fact, \(\frac{v}{d}\) is an odd number and \(d \neq 3\). Indeed if \(\frac{v}{d} = 3\) we would have \(t = d\) and hence \(a + b \geq \frac{d-1}{2} = \frac{2d-1}{2}\) but we are assuming \(a + b < \frac{t-r-1}{2}\). Note that, for any \(i\) and \(j\), \(m_{i,j+1} - m_{i,j} = t \pmod{v}\), namely two elements adjacent on a row form an edge of length \(t\). Also \(|m_{i+1,j} - m_{i,j}| = 2\) if \(i\) is odd, while \(|m_{i+1,j} - m_{i,j}| = 1\) if \(i\) is even. Our aim is to construct a edges of length 1 and \(b\) edges of length 2 in such a way that the elements of \(M\) not used to obtain these edges appear in all but one rows (say \(R\)) as an even number of adjacent entries. In fact, if this holds, then it will be immediate to construct the edges of length \(t\) with pairs of elements adjacent on a row. Obviously, the isolated vertex of the near 1-factor will belong to the row \(R\). So, it suffices to explain how to construct the edges of length 1 and 2. Then, the reader can easily check that the previous condition is always satisfied in the following constructions. We have to split the proof into two cases.
Case 1. \(a + b - \frac{d-1}{2}\) even.

By the assumptions on \(a, b\) and \(d\) it is possible to find two positive integers \(\tilde{a}\) and \(\tilde{b}\) such that \(\tilde{a} + \tilde{b} = \frac{d-1}{2}\), \(a - \tilde{a} = 2\alpha\) and \(b - \tilde{b} = 2\beta\), for suitable \(\alpha\) and \(\beta\). Clearly, in general, it is possible more than one choice for the pair \((\tilde{a}, \tilde{b})\). By Corollary 13 there exists a near 1-factor \(F\) of \(K_d\) such that \(\delta(F) = \{1^5, 2^5\}\). We start taking the edges of \(F\). Now we have to construct \(2\alpha\) edges of length 1 and \(2\beta\) edges of length 2. Let \(2\alpha = (\frac{\tilde{q}}{2} - 1)\tilde{q} + \tilde{r}\) with \(0 \leq \tilde{r} < \frac{\tilde{q}}{2} - 1\), note that \(\tilde{r}\) is even. Take the following \(2\alpha\) edges of length 1: \([m_{2i,j}, m_{2i+1,j}]\) for \(i = \frac{d-1}{2} - \tilde{q} + 1, \ldots, \frac{d-1}{2}\), \(j = 2, \ldots, \frac{\tilde{q}}{2}\) (if \(\tilde{q} > 0\)) and \([m_{2\tilde{q}+1,j}, m_{2\tilde{q}+2,j}]\) for \(j = 2, \ldots, \frac{\tilde{q}}{2} + 1\). Let \(2\beta = (\frac{\tilde{q}}{2} - 1)\tilde{q} + \tilde{r}\) with \(0 \leq \tilde{r} < \frac{\tilde{q}}{2} - 1\), note that \(\tilde{r}\) is even. Take the following \(2\beta\) edges of length 2: \([m_{2i+1,j}, m_{2i+2,j}]\) for \(i = 0, \ldots, \tilde{q} - 1, j = 2, \ldots, \frac{\tilde{q}}{2}\) (if \(\tilde{q} > 0\)) and \([m_{2\tilde{q}+1,j}, m_{2\tilde{q}+2,j}]\) for \(j = \frac{\tilde{q}}{2} - \tilde{r} + 1, \ldots, \tilde{q}\).

Since \(c \geq \frac{\tilde{q}^2 - \frac{d}{2}}{2}\), the edges so constructed involve distinct vertices of \(K_v\). Also, since \(\tilde{r}\) and \(\tilde{r}\) are even one can check that the elements of \(M\) not used to obtain these edges appear in all but one rows as an even number of adjacent entries.

Case 2. \(a + b - \frac{d-1}{2}\) odd.

By the assumptions on \(a, b\) and \(d\) it is possible to find two positive integers \(\tilde{a}\) and \(\tilde{b}\) such that \(\tilde{a} + \tilde{b} = \frac{d-1}{2}\), \(a - \tilde{a} = 2\alpha + 1\) and \(b - \tilde{b} = 2\beta\), for suitable \(\alpha\) and \(\beta\). Also in this case, it may be possible more than one choice for the pair \((\tilde{a}, \tilde{b})\). From Corollary 13 there exists a near 1-factor \(F\) of \(K_d\) such that \(\delta(F) = \{1^5, 2^5\}\) whose isolated vertex is in \(\{d - 2, d - 1\}\), namely in \(\{m_{d-1,1}, m_{d,1}\}\). We start taking the edges of \(F\). Next, we take the edges of length 2. Let \(2\beta = (\frac{\tilde{q}}{2} - 1)\tilde{q} + \tilde{r}\) with \(0 \leq \tilde{r} < \frac{\tilde{q}}{2} - 1\), note that \(\tilde{r}\) is even. We consider the edges \([m_{2i+1,j}, m_{2i+2,j}]\), for \(i = 0, \ldots, \tilde{q} - 1\) and \(j = 2, \ldots, \frac{\tilde{q}}{2}\) (if \(\tilde{q} > 0\)), and the edges \([m_{2\tilde{q}+1,j}, m_{2\tilde{q}+2,j}]\), for \(j = \frac{\tilde{q}}{2} - \tilde{r} + 1, \ldots, \tilde{q}\).

Now, we take the edges of length 1. Firstly we take the edge \([m_{d-1,3}, m_{d,3}]\). Next we distinguish two cases. If \(2\alpha \leq \frac{\tilde{q}}{2} - 3\) we take the edges \([m_{d-1,j}, m_{d,j}]\) for \(j = 4, \ldots, 2\alpha + 3\). If \(2\alpha > \frac{\tilde{q}}{2} - 3\), let \(2\alpha - (\frac{\tilde{q}}{2} - 3) = (\frac{\tilde{q}}{2} - 1)\tilde{q} + \tilde{r}\) with \(0 \leq \tilde{r} < \frac{\tilde{q}}{2} - 1\), note that \(\tilde{r}\) is even. We consider the edges \([m_{d-1,j}, m_{d,j}]\) for \(j = 4, \ldots, \frac{\tilde{q}}{2}\), \([m_{d-2,j}, m_{d-2-1,j}]\), for \(i = 1, \ldots, \tilde{q}\) and \(j = 2, \ldots, \frac{\tilde{q}}{2}\) (if \(\tilde{q} > 0\)), and \([m_{d-2\tilde{q}-2,j}, m_{d-2\tilde{q}-3,j}]\), for \(j = 2, \ldots, \tilde{r} + 1\). Since \(c \geq \frac{\tilde{q}^2 - d}{2}\), the edges so constructed involved distinct vertices of \(K_v\). The isolated vertex is in \(\{m_{d-1,2}, m_{d,2}\}\).

**Example 19.** Let \(L = \{1^3, 2^2, 21^{19}\}\). Hence, we have \(v = 49\), \(d = 7\) and \(a + b - \frac{d-1}{2} = 3 + 2 - 3 = 2\) is even. So we are in Case 1 of Proposition 18. We start constructing the \(7 \times 7\) matrix \(M\):

\[
M = \begin{pmatrix}
\emptyset & 21 & 42 & 14 & 35 & 7 & 28 \\
2 & 23 & 44 & 16 & 37 & 9 & 30 \\
1 & 22 & 43 & 15 & 36 & 8 & 29 \\
3 & 24 & 45 & 17 & 38 & 10 & 31 \\
4 & 25 & 46 & 18 & 39 & 11 & 32 \\
6 & 27 & 48 & 20 & 41 & 13 & 34 \\
5 & 26 & 47 & 19 & 40 & 12 & 33 \\
\end{pmatrix}.
\]

Note that in this case we have more than one choice for \(\tilde{a}\) and \(\tilde{b}\), in fact we can choose \((\tilde{a}, \tilde{b}) = (3, 0)\) or \((1, 2)\). In this example we choose \(\tilde{a} = 3\) and \(\tilde{b} = 0\). So we take \(F = \{[2, 1], [3, 4], [6, 5]\} \cup \{0\}\). Now we have \(a - \tilde{a} = 0\) and \(b - \tilde{b} = 2\), hence \(\alpha = 0\) and \(\beta = 1\). So we have to take the following edges of length 2: \([7, 9]\)
Example 20. Let \(\gcd(t, v) = 2\). Assume that \(\alpha = 2\) and \(\beta = 1\). So, we have to take the following edges of length 2: \(\{0, 2\}\), \(\{0, 1\}\), and \(\{0, 0\}\). Take also \(m_{4,3} = 53\). Clearly, the isolated vertex is 0.

Example 21. Let \(L = \{1^7, 2^4, 25^{16}\}\). Hence, we have \(v = 55\), \(d = 5\) and \(a + b - 2 = 9\) odd. So we are in Case 2 of Proposition 18. We start constructing the \(5 \times 11\) matrix \(M\):

\[
M = \begin{pmatrix}
0 & 25 & 50 & 20 & 45 & 15 & 40 & 10 & 35 & 5 & 30 \\
2 & 27 & 52 & 22 & 47 & 17 & 42 & 12 & 37 & 7 & 32 \\
1 & 26 & 51 & 21 & 46 & 16 & 41 & 11 & 36 & 6 & 31 \\
3 & 28 & 53 & 23 & 48 & 18 & 43 & 13 & 38 & 8 & 33 \\
4 & 29 & 54 & 24 & 49 & 19 & 44 & 14 & 39 & 9 & 34
\end{pmatrix}.
\]

Note that in this case we have more than one choice for \(\tilde{a}\) and \(\tilde{b}\), in fact we can choose \((\tilde{a}, \tilde{b}) = (0, 2)\) or \((2, 0)\). In this example we choose \(\tilde{a} = 0\) and \(\tilde{b} = 2\). So we take \(F = \{\{0, 2\}, \{1, 3\}\} \cup \{4\}\). Now we have \(a - \tilde{a} = 7\) and \(b - \tilde{b} = 2\), hence \(\alpha = 3\) and \(\beta = 1\). So we have to take the following edges of length 2: \(\{5, 7\}\) and \(\{30, 32\}\). Take also \(m_{4,3} = 53\). Also, take the following edges of length 1: \(\{23, 24\}, \{48, 49\}, \{18, 19\} \cup \{43, 44\}, \{13, 14\}, \{38, 39\}\). The elements used up to now to construct the edges of length 1 and 2 are highlighted in bold in the matrix. Finally, it is easy to construct 16 edges of length 25 as follows: \(\{25, 50\}, \{20, 45\}, \{15, 40\}, \{10, 35\}, \{27, 52\}, \{22, 47\}, \{17, 42\}, \{12, 37\}, \{26, 51\}, \{21, 46\}, \{16, 41\}, \{11, 36\}, \{6, 31\}, \{8, 33\}, \{4, 29\}, \{9, 34\}\). Clearly, the isolated vertex is 28.

In order to complete the study of \(\text{MPP}([1^a, 2^b, t^c])\) we are left to consider the case \(\gcd(t, v) = 1\) and \(a + b < \left\lfloor \frac{t-1}{2} \right\rfloor\), where \(v = 2(a + b + c) + 1\). In the following we give some partial results about this case. So we take \(a + b < \left\lfloor \frac{t-1}{2} \right\rfloor\), but not “too small”. Clearly, we believe that \(\text{MPP} \) conjecture holds also for \(a + b \) “very small” and we will give some examples for these cases at the end of this section.

Proposition 21. Let \(L = \{1^a, 2^b, t^c\}\), where \(t\) is an integer with \(t > 2\). Let \(v = 2(a + b + c) + 1\) and let \(r\) be the remainder of the division of \(v\) by \(2t\). Assume that \(a + b \geq t - \frac{t+1}{2}\) if \(\frac{v}{2t}\) is even and that \(a + b \geq t - \frac{t-1}{2}\) if \(\frac{v}{2t}\) is odd. Then there exists a near 1-factor \(F\) of \(K_v\) such that \(t(F) = L\).

Proof. In view of Propositions 16 and 18, we may assume \(a + b < \left\lfloor \frac{t-1}{2} \right\rfloor\) and \(\gcd(t, v) = 1\). Let \(v = qt + r\), with \(0 < r < t\). We construct an incomplete matrix \(M\) with \(t\) rows and \(q + 1\) columns whose elements are \(\{0, 1, 2, \ldots, v - 1\}\), namely the vertices of \(K_v\), and the element on the \(i\)-th row and \(j\)-th column is \(m_{i,j} = (i - 1) + (j - 1)t\), for \(i = 1, \ldots, t\), \(j = 1, \ldots, q\) and \(i = 1, \ldots, r\) if \(j = q + 1\). Hence the first \(q\) columns are complete, while in the last column we have only \(r\) elements. It is easy to see that \(m_{i,j+1} - m_{i,j} = t\) for any \(i, j\), hence two elements adjacent on a row can be connected by an edge of length \(t\). Also, we have \(m_{i+1,j} - m_{i,j} = 1\) for any \(i, j\), so two elements adjacent on a column can be connected by an edge of length 1.

As in Proposition 18, our aim is to construct \(a\) edges of length 1 and \(b\) edges of length 2 in such a way that the elements of \(M\) not used to obtain these edges appear in all but one rows as an even number of adjacent entries. So, in the following, we
will explain how to construct the edges of length 1 and 2 in such a way that the previous condition is always satisfied.

Since $v$ is odd and $r'$ is the remainder of the division of $v$ by $2t$, $r'$ is an odd integer too. Now we split the construction into two cases.

Case 1. $q$ even.
Note that, in this case, the rows with an odd number of elements are the first $r' = r$ rows of $M$.

Case 1A. $a + b = \frac{r' + 1}{2}$.
By Corollary 13, there exists a near 1-factor $F$ of $K_{r'}$ with $\delta(F) = \{1^a, 2^b\}$. Then it suffices to take the edges of $F$.

Case 1B. $a + b > \frac{r' + 1}{2}$.
Firstly, consider the case $r' = 1$. We apply Corollary 13 to obtain a near 1-factor $F$ such that $\delta(F) = \{1^a, 2^b\}$ and so it is always satisfied.

By Corollary 13, there exists a near 1-factor $F$ of $K_{r'}$ with $\delta(F) = \{1^a, 2^b\}$. Then it suffices to take the edges of $F$.

Firstly, consider the case $r' = 1$. We apply Corollary 13 to obtain a near 1-factor $F$ such that $\delta(F) = \{1^a, 2^b\}$ and so it is always satisfied.

Next, in $K_v$ we take the edges of the graphs $F + t - 2|L'| - 1$ and $F + 2t - 2|L'| - 1$. If $a$ and $b$ are both even, we have done. If $a$ is even and $b$ is odd, we still have to construct an edge of length 2, take $[1, v - 1]$. If $a$ is odd and $b$ is even, we still have to construct an edge of length 1, take $[0, v - 1]$. Finally, if $a$ and $b$ are both odd, we still have to construct an edge of length 1 and one of length 2, take $[v - t - 1, v - t]$ and $[1, v - 1]$.

Assume now $r' > 1$. The construction in this subcase depends on the parity of $a + b - \frac{r' + 1}{2}$. It is easy to see that if $a + b - \frac{r' + 1}{2}$ is even there exist $\tilde{a}$, $\tilde{b}$ such that $\tilde{a} + \tilde{b} = \frac{r' + 1}{2}$ and $a - \tilde{a} = 2\alpha$ and $b - \tilde{b} = 2\beta$, for suitable $\alpha$ and $\beta$. Now, it is possible, in view of Corollary 13, to construct a near 1-factor $F$ of $K_{r'}$ such that $\delta(F) = \{1^a, 2^b\}$ and a near 1-factor $F'$ of $K_{2(\alpha + \beta) + 1}$ such that $\delta(F') = \{1^\alpha, 2^\beta\}$.

Note that, since $a + b < \frac{r' + 1}{2}$, we have $2(\alpha + \beta) + 1 < t$. Now, in $K_v$ we take the edges of $F + tq$, $F' + t - 2(\alpha + \beta) - 1$ and $F' + 2t - 2(\alpha + \beta) - 1$, these are exactly $a$ edges of length 1 and $b$ edges of length 2.

It is easy to see that if $a + b - \frac{r' + 1}{2}$ is odd, there exist $\tilde{a}$, $\tilde{b}$ such that $\tilde{a} + \tilde{b} = \frac{r' + 1}{2}$ and $a - \tilde{a} = 2\alpha + 1$ and $b - \tilde{b} = 2\beta$, for suitable $\alpha$ and $\beta$. Now, in $K_v$ we take again the edges of $F + tq$, $F' + t - 2(\alpha + \beta) - 1$ and $F' + 2t - 2(\alpha + \beta) - 1$, where $F$ and $F'$ are defined as before. Furthermore, we take also the edge $[m_{i, 1}, m_{i, \tilde{a} + 1}]$ of length 1, where $i$ is the isolated vertex of $F$ (observe that we may always assume $i \geq 1$). Finally, observe that, since $r' < t$ and $a + b < \frac{r' + 1}{2}$ it follows $t - 2(\alpha + \beta) - 1 > r'$ and so $F + tq$ and $F' + t - 2(\alpha + \beta) - 1$ involve distinct rows.

Case 2. $q$ odd.
Note that, in this case, the rows with an odd number of elements are the last $2t - r' = t - r$ rows of $M$.

Case 2A. $a + b = t - \frac{r' + 1}{2}$
By Corollary 13 there exists a near 1-factor $F$ such that $\delta(F) = \{1^a, 2^b\}$ so it is sufficient to take the edges of $F + t - 2(a + b) - 1$ in $K_v$.

Case 2B. $a + b > t - \frac{r' + 1}{2}$.
If $a + b - (t - \frac{r' + 1}{2})$ is even there exist $\tilde{a}$ and $\tilde{b}$ such that $\tilde{a} + \tilde{b} = t - \frac{r' + 1}{2}$ and $a - \tilde{a} = 2\alpha$ and $b - \tilde{b} = 2\beta$, for suitable $\alpha$ and $\beta$. Now let $F$ and $F'$ as in Case 1B. In $K_v$ we take the edges of $F + v - t$, $F'$ and $F' + t$.

If $a + b - (t - \frac{r' + 1}{2})$ is odd there exist $\tilde{a}$ and $\tilde{b}$ such that $\tilde{a} + \tilde{b} = t - \frac{r' + 1}{2}$ and $a - \tilde{a} = 2\alpha + 1$ and $b - \tilde{b} = 2\beta$, for suitable $\alpha$ and $\beta$. In $K_v$ we take again the edges
Finally, we take the edge \( m \) Corollary 23. v = 16 = 0, so we consider \( \beta \) 21. Consider the edges \( m \) For thesis follows from Propositions 16, 18 and 21. \[
M = \begin{pmatrix}
0 & 12 & 24 \\
1 & 13 & 25 \\
2 & 14 & 26 \\
3 & 15 & 27 \\
4 & 16 & 28 \\
5 & 17 & \\
6 & 18 & \\
7 & 19 & \\
8 & 20 & \\
9 & 21 & \\
10 & 22 & \\
11 & 23 &
\end{pmatrix}
\quad M' = \begin{pmatrix}
0 & 12 & 24 & 36 \\
1 & 13 & 25 & 37 \\
2 & 14 & 26 & 38 \\
3 & 15 & 27 & 39 \\
4 & 16 & 28 & 40 \\
5 & 17 & 29 & \\
6 & 18 & 30 & \\
7 & 19 & 31 & \\
8 & 20 & 32 & \\
9 & 21 & 33 & \\
10 & 22 & 34 & \\
11 & 23 & 35 &
\end{pmatrix}
\]

**Figure 1.** Matrices of Example 22.

of \( F + v - t, F' \) and \( F' + t \). Finally we take the edge \([m_{i,1}, m_{i+1,1}]\) where \( i \) is the isolated vertex of \( F + r \).

\[\Box\]

**Example 22.** Let \( L = \{1^2, 2^3, 12^9\} \), hence, following the notation of Proposition 21, we have \( v = 29, t = 12, q = 2 \) and \( r = r' = 5 \). It is easy to see that we are in Case 1B of the proof of Proposition 21. Let \( M \) be the incomplete matrix defined in Proposition 21 (see Fig. 1). We take \( \tilde{a} = \tilde{b} = 1 \) and \( F = \{[0, 1], [2, 4]\} \cup \{3\} \). Now \( a = 0 \) and \( \beta = 1 \) and we consider \( F' = \{[0, 2]\} \cup \{1\} \). Following the proof, we have to take the edges of \( F + 24, F' + 9 \) and \( F' + 21 \), namely the following edges \([24, 25], [26, 28], [9, 11], [21, 23]\). Since the isolated vertex of \( F \) is 3 we have to take also the edge \([m_{3,1}, m_{4,1}] = [2, 3]\). Up to now we have constructed all the edges of length 1 and 2. It is easy to see that the elements not used to construct these edges appear in an even number on all but one row. Hence it is immediate to construct the edges of length 12.

Now let \( L = \{1^4, 2^3, 12^{14}\} \), hence we have \( v = 41, t = 12, q = 3, r = 5 \) and \( r' = 17 \). So we are in Case 2B of the proof of Proposition 21. Let \( M' \) be the incomplete matrix defined in Proposition 21 (see Fig. 1). Since \( a + b - (t - \frac{r}{2} - \frac{r'}{2}) \) is odd, we take \( \tilde{a} = 1, \tilde{b} = 2 \) and we consider \( F = \{[0, 2], [1, 3], [4, 5]\} \cup \{6\} \). Now \( a = 1 \) and \( \beta = 0 \), so we consider \( F' = \{[0, 1]\} \cup \{2\} \). Following the proof now we have to take the edges of \( F + 29, F' \) and \( F' + 12 \), namely \([29, 31], [30, 32], [33, 34], [0, 1], [12, 13]\). Finally, we take the edge \([m_{11,1}, m_{12,1}] = [10, 11]\) since the isolated vertex of \( F + r \) is 11. Now it is trivial to construct the edges of length 12.

**Corollary 23.** If \( t \leq 11 \), \( \text{MPP}(\{1^a, 2^b, t^c\}) \) holds for any \( a, b, c \geq 0 \).

**Proof.** For \( t = 3, 4, 5, 6 \) the thesis directly follows from Proposition 16. If \( t = 7 \) the thesis follows from Propositions 16, 18 and 21.

Take now \( t = 8 \), in view of Propositions 16 and 21 we are left to consider the following classes of lists \( \{1, 2, 8^{q} + 1\} \) and \( \{1, 2, 8^{q} + 2\} \). If \( L = \{1, 2, 8^{q} + 1\} \), \( v = 16q + 7 \). Let \( q = 2q' \) and \( M \) be the incomplete matrix defined in Proposition 21. Consider the edges \([m_{2,1}, m_{1,q+1}]\) of length 8, \([m_{3,1}, m_{4,1}]\) of length 1 and \([m_{5,1}, m_{7,1}]\) of length 2. Now all the rows of \( M \) except the 6-th one have an even number of adjacent entries so we can construct the remaining edges of length 8.
If \( L = \{1, 2, 8q^2+2\} \), \( v = 16q' + 9 \). Let \( q = 2q' + 1 \) and \( M \) be the incomplete matrix defined in Proposition 21. Consider the edges \([m_{2,1}, m_{3,q}]\) of length 8, \([m_{4,1}, m_{5,1}]\) of length 1 and \([m_{6,1}, m_{8,1}]\) of length 2. Now all the rows of \( M \) except the 7-th one have an even number of adjacent entries so we can construct the remaining edges of length 8.

For \( t = 9 \), in view of Propositions 16, 18 and 21 we are left to consider the following classes of lists \( \{1, 2, 9q' + 1\} \) and \( \{1, 2, 9q' + 3\} \). If \( L = \{1, 2, 9q' + 1\} \), \( v = 18q' + 7 \). Let \( q = 2q' \) and \( M \) be the incomplete matrix defined in Proposition 21. Consider the edges \([m_{3,1}, m_{1, q+1}]\) of length 9, \([m_{2,1}, m_{4,1}]\) of length 2 and \([m_{5,1}, m_{6,1}]\) of length 1. Now all the rows of \( M \) except the 7-th one have an even number of adjacent entries so we can construct the remaining edges of length 9. If \( L = \{1, 2, 9q' + 3\} \), \( v = 18q' + 11 \). Let \( q = 2q' + 1 \) and \( M \) be the incomplete matrix defined in Proposition 21. Consider the edges \([m_{7,1}, m_{9,q}]\) of length 9, \([m_{3,1}, m_{4,1}]\) of length 1 and \([m_{6,1}, m_{8,1}]\) of length 2. Now all the rows of \( M \) except the 5-th one have an even number of adjacent entries so we can construct the remaining edges of length 9.

Finally, for \( t = 10, 11 \) it is easy to see that we are left to consider eight classes of lists for each value of \( t \). By direct computation we have checked that the conjecture holds also in these cases.

To conclude we present an example for lists not considered in previous propositions. We consider the “most difficult” case namely when \( a = b = 1 \). It is important to underline that the strategy used for these particular cases can be generalized to infinite classes of lists.

**Example 24.** Take \( L = \{1, 2, 19^{23}\} \), hence \( v = 51 \). Let \( M \) be the incomplete matrix defined in Proposition 21 (see Fig. 2) and take the following edges \([38, 6], [39, 7], [40, 8], [41, 9]\) of length 19, \([42, 43]\) of length 1 and \([48, 50]\) of length 2. Now all the rows of \( M \) except the 12-th one have an even number of adjacent entries so we can construct the remaining edges of length 19. It is easy to see that this construction can be easily generalized to any list \( \{1, 2, 19^{19k+4}\} \).

Take now \( L = \{1, 2, 19^{28}\} \), hence \( v = 61 \). Let \( M' \) be the incomplete matrix defined in Proposition 21 (see Fig. 2) and take the following edges \([57, 15], [58, 16], [59, 17], [60, 18], [0, 42], [1, 43], [2, 44], [3, 45], [8, 50]\) of length 19, \([9, 11]\) of length 2 and \([13, 14]\) of length 1. Now all the rows of \( M' \) except the 11-th one have an even number of adjacent entries so we can construct the remaining edges of length 19. It is easy to see that this construction can be easily generalized to any list \( \{1, 2, 19^{19k+4}\} \).

We have to point out that the constructions illustrated in the previous example work for any values of \( t \) if \( L = \{1, 2, t^c\} \) with \( c \) belonging to some suitable congruence classes modulo \( t \). Indeed for \( t = 19 \) we are able to prove MPP(L) for any \( c \), but the construction of the near 1-factor depends on the congruence class of \( c \) modulo 19. In our opinion it seems not possible to present a general construction, but it is necessary to split the proof of the remaining open case in several subcases.

**Acknowledgments**

The authors thank the anonymous referees for their useful suggestions and comments.
16 ANITA PASOTTI AND MARCO ANTONIO PELLEGRINI

\[
M = \begin{pmatrix}
0 & 19 & 38 \\
1 & 20 & 39 \\
2 & 21 & 40 \\
3 & 22 & 41 \\
4 & 23 & 42 \\
5 & 24 & 43 \\
6 & 25 & 44 \\
7 & 26 & 45 \\
8 & 27 & 46 \\
9 & 28 & 47 \\
10 & 29 & 48 \\
11 & 30 & 49 \\
12 & 31 & 50 \\
13 & 32 & \\
14 & 33 & \\
15 & 34 & \\
16 & 35 & \\
17 & 36 & \\
18 & 37 & \\
\end{pmatrix}
\]

\[
M' = \begin{pmatrix}
0 & 19 & 38 & 57 \\
1 & 20 & 39 & 58 \\
2 & 21 & 40 & 59 \\
3 & 22 & 41 & 60 \\
4 & 23 & 42 & \\
5 & 24 & 43 & \\
6 & 25 & 44 & \\
7 & 26 & 45 & \\
8 & 27 & 46 & \\
9 & 28 & 47 & \\
10 & 29 & 48 & \\
11 & 30 & 49 & \\
12 & 31 & 50 & \\
13 & 32 & \\
14 & 33 & \\
15 & 34 & \\
16 & 35 & \\
17 & 36 & \\
18 & 37 & 56 & \\
\end{pmatrix}
\]

Figure 2. Matrices of Example 24.

REFERENCES

[1] Baker C.A., Extended Skolem sequences. J. Combin. Des. 3, 363–379 (1995).
[2] Bonvicini S., Buratti M., Rinaldi G., Traetta T., Some progress on the existence of 1-rotational Steiner triple systems. Des. Codes Cryptogr. 62 (2012), 63–78.
[3] Bryant D., El-Zanati S., Graph decompositions, In: C.J. Colbourn and J.H. Dinitz editors, Handbook of combinatorial designs, Second Edition, Chapman & Hall/CRC, Boca Raton, FL, 2006, pp. 477–486.
[4] Buratti M., 1-rotational Steiner triple systems over arbitrary groups. J. Combin. Des. 9, 215–226 (2001).
[5] Buratti M., Merola F., Dihedral Hamiltonian Cycle Systems of the Cocktail Party Graph. J. Combin. Des. 21, 1–23 (2013).
[6] Buratti M., Merola F., Hamiltonian cycle systems which are both cyclic and symmetric. J. Combin. Des. 22, 367–390 (2014).
[7] Capparelli S., Del Fra A., Hamiltonian paths in the complete graph with edge-lengths 1, 2, 3. Electron. J. Combin. 17 (2010), gR44.
[8] Dinitz J.H., Starters, In: Colbourn C.J. and Dinitz J.H. editors, Handbook of Combinatorial Designs, Second Edition, Chapman & Hall/CRC, Boca Raton, FL, 2006, pp. 622–628.
[9] Dinitz J.H., Janiszewski S.R., On hamiltonian Paths with Prescribed edge lengths in the Complete Graph. Bull. Inst. Combin. Appl. 57, 42–52 (2009).
[10] Francetic N., Mendelsohn E., A survey of Skolem-type sequences and Rosa’s use of them. Math. Slovaca 59, 39–76 (2009).
[11] Godsil C., Royle G., Algebraic graph theory. Graduate Texts in Mathematics. Vol 207. Springer, (2001).
[12] Horak P., Rosa A., On a problem of Marco Buratti. Electron. J. Combin. 16 (2009), gR20.
[13] Pasotti A., Pellegrini M.A., A new result on the problem of Buratti, Horak and Rosa. Discrete Math. 319, 1–14 (2014).
[14] Pasotti A., Pellegrini M.A., On the Buratti-Horak-Rosa conjecture about Hamiltonian paths in complete graphs. Electron. J. Combi. 21 (2014), gP2.30.
[15] Rosa A., On a problem of Mariusz Meszka. Discrete Math. 338, 139–143 (2015).
[16] Shalaby N., Skolem and Langford Sequences, In: C.J. Colbourn and J.H. Dinitz editors, Handbook of combinatorial designs, Second Edition, Chapman & Hall/CRC, Boca Raton, FL, 2006, pp. 612–616.
[17] West D., Introduction to graph theory. Prentice Hall, New Jersey (1996).
[18] West D., http://www.math.uiuc.edu/~west/regs/buratti.html. Accessed 28 March 2014.
[19] Wu S.-L., Buratti M., A complete solution to the existence problem for 1-rotational \( k \)-cycle systems of \( K_v \). J. Combin. Des. 17 (2009), 283–293.

DICATAM - SEZ. MATEMATICA, UNIVERSITÀ DEGLI STUDI DI BRESCIA, VIA BRANZE 43, I-25123 BRESCIA, ITALY

E-mail address: anita.pasotti@unibs.it

DIPARTIMENTO DI MATEMATICA E FISICA, UNIVERSITÀ CATTOLICA DEL SACRO CUORE, VIA MUSEI 41, I-25121 BRESCIA, ITALY

E-mail address: marcoantonio.pellegrini@unicatt.it