Braided bialgebras in a generated monoidal Ab-category

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Abstract

We start from any small strict monoidal braided Ab-category and extend it to a monoidal nonstrict braided Ab-category which contains braided bialgebras. The objects of the original category turn out to be modules for these bialgebras.

0 Introduction

The notion of bialgebras and Hopf algebras in braided categories were introduced by S. Majid in [4]. He has considered a braided monoidal (tensor) category and in the usual definitions of algebras, coalgebras, bialgebras and Hopf algebras, he has replaced the flip by the braiding in the obvious way. Majid has called braided bialgebra to a bialgebra in a braided category. Refer to [1] and [2] for generalities in braided monoidal categories and to [4], [5] and [6] for the definition and results in the theory of braided bialgebras and braided Hopf algebras.

The purpose of this paper is to present a construction in which starting from a small braided monoidal Ab-category \( \mathcal{C} \) and an infinite set \( S_0 \) we create a new monoidal braided category \( \mathcal{C}^{S_0} \) that contains the original \( \mathcal{C} \) as a subcategory, and more important, it contains objects with bialgebra structure, in such a way that the objects of the original category \( \mathcal{C} \) are modules for these bialgebras. Remember that a category \( \mathcal{C} \) is said to be an Ab-category (also called pre-abelian category; cf. [7]) if for any pair of objects \( V, W \) its set of morphisms \( \text{hom}(V, W) \) is an additive abelian group and the composition of morphisms is bilinear. In the context of monoidal Ab-categories we shall assume that the tensor product of morphisms is bilinear. For the construction we proceed as follows. Section 1 is divided into two parts; in the first part, out of any small Ab-category \( \mathcal{C} \) and any set \( S_0 \), we construct the new category \( \mathcal{C}^{S_0} \) which is also an Ab-category. In the second part we assume that \( \mathcal{C} \) is strict monoidal and that \( S_0 \) is infinite, and so, we extend the monoidal structure to \( \mathcal{C}^{S_0} \). However, the extended monoidal structure is not strict, so we have to work with associative constraints, and left and right units. In the second part we show how to extend a braiding and a twist from \( \mathcal{C} \) to \( \mathcal{C}^{S_0} \). Since the new category is nonstrict monoidal, we need to define algebras, coalgebras, bialgebras and modules in this case. This is easily made
if in the categorical definitions of the latter notions we substitute the equalities by an equivalent relation in the set of morphisms of $\mathcal{C}^{S_0}$. Roughly speaking, we declare two morphisms of $\mathcal{C}^{S_0}$ related if their domains and codomains are related by associativity and/or units. This relation, obviously, agrees with the identity if the category is strict; this is explained in detail at the end of section 1. We start section 2 defining algebras, coalgebras, bialgebras and modules in nonstrict monoidal categories in general, and then state and prove the main theorem (Theorem 2.1) of this paper. Throughout the proof we use graphical calculus as explained in [1] and [2]. This work is influenced by Yetter’s paper [3].

1 The category $\mathcal{C}^{S_0}$

Let $\mathcal{C}$ be a small Ab-category. We shall denote by $\text{Obj}(\mathcal{C})$ and $\mathcal{H}$ the sets of its objects and morphisms, respectively. We are going to associate to $\mathcal{C}$ a new category $\mathcal{C}^{S_0}$ as follows. Let us take a fixed set $S_0$ and consider the set $\mathcal{M}(S_0, \text{Obj}(\mathcal{C})) = \{ S_0 \supset S_f \overset{f}{\longrightarrow} \text{Obj}(\mathcal{C}) \}$, where $S_f$ is any subset of $S_0$ and $f$ is a set-theoretical function. The objects of $\mathcal{C}^{S_0}$ will be the elements of $\mathcal{M}(S_0, \text{Obj}(\mathcal{C}))$. Let $f : S_f \longrightarrow \text{Obj}(\mathcal{C})$ and $g : S_g \longrightarrow \text{Obj}(\mathcal{C})$ be two objects. A morphism $F : f \longrightarrow g$ will be a two-variable function $F : S_f \times S_g \longrightarrow \mathcal{H}$ such that:

(i) $F(x, y) : f(x) \longrightarrow g(y)$, for all $(x, y) \in S_f \times S_g$.

(ii) If $S_g$ is infinite, then for each $x \in S_f$ there exists a finite set $S^F_x \subset S_g$, such that $F(x, y) = 0$ if $y \in S_g - S^F_x$.

Let $f : S_f \longrightarrow \text{Obj}(\mathcal{C})$, $g : S_g \longrightarrow \text{Obj}(\mathcal{C})$, and $h : S_h \longrightarrow \text{Obj}(\mathcal{C})$ be two objects, and $F : f \longrightarrow g$, $G : g \longrightarrow h$ be morphisms. Define $G \circ F : f \longrightarrow h$ as the function $G \circ F : S_f \times S_h \longrightarrow \mathcal{H}$ given by:

$$
(G \circ F)(x, y) = \sum_{z \in S^F_g} G(z, y) \circ F(x, z) \tag{1}
$$

for $x \in S_f$ and $y \in S_h$. This sum is always finite. Indeed, if we write $S^F_x = \{ z_1, \ldots, z_k \}$, then the sum becomes

$$
(G \circ F)(x, y) = \sum_{i=1}^{k} G(z_i, y) \circ F(x, z_i) \tag{2}
$$

It is clear that the function $G \circ F$ satisfies condition (i). Besides, if $y \notin S^G_{z_1} \cup \ldots \cup S^G_{z_k}$, then $G(z_i, y) = 0$ for $1 \leq i \leq k$, so if we choose $S^G_x = S^G_{z_1} \cup \ldots \cup S^G_{z_k}$, then we have $y \in S_h - S^G_x$, thus $(G \circ F)(x, y) = 0$. Therefore $G \circ F$ also satisfies condition (ii).
For any \( f : S_f \to \text{Obj}(C) \) define \( \text{Id}_f : f \to f \) as the function \( \text{Id}_f : S_f \times S_f \to \mathcal{H} \), given by \( \text{Id}_f(x, y) = \delta_{x,y} \text{id}_f(x) : f(x) \to f(y) \) for \((x, y) \in S_f \times S_f\).

For \( G : f \to g \) one has

\[
(G \circ \text{Id}_f)(x, y) = \sum_{z \in S_f} G(z, y) \circ \text{Id}(x, z) \tag{3}
\]

Therefore \( G \circ \text{Id}_f = G \). Analogously \( \text{Id}_g \circ G = G \) for any morphism \( G : f \to g \).

Furthermore this operation is associative. Indeed, if \( F : f \to g \), \( G : g \to h \), and \( H : h \to i \), then

\[
((H \circ G) \circ F)(w, z) = \sum_{x \in S_g} (H \circ G)(x, z) \circ F(w, x) \\
= \sum_{x \in S_g} \sum_{y \in S_h} (H(y, z) \circ G(x, y)) \circ F(w, x) \\
= \sum_{y \in S_h} H(y, z) \circ (\sum_{x \in S_g} G(x, y) \circ F(w, x)) \tag{4}
\]

Hence we have proved that \( \mathcal{C}^{S_0} \) is a category. If for two morphisms \( F, G : f \to g \) we define the function \( (F + G)(x, y) = F(x, y) + G(x, y) \), which trivially satisfies conditions (i) and (ii), we see that \( \mathcal{C}^{S_0} \) is also an Ab-category. The following proposition proves that the direct sum of certain collections of objects in \( \mathcal{C}^{S_0} \) is defined.

**Proposition 1.1.** Let \( \{f_i : S_i \to \text{Obj}(C)\}_{i \in I} \) be any collection of functions such that the sets \( S_i, i \in I \), are pairwise disjoint subsets of \( S_0 \). Then \( \{f_i : S_i \to \text{Obj}(C)\}_{i \in I} \) is the coproduct of \( \{f_i : S_i \to \text{Obj}(C)\}_{i \in I} \) in \( \mathcal{C}^{S_0} \).

**Proof.** Suppose we are given an object \( g : S_g \to \text{Obj}(C) \) and a family of morphisms \( T_i : (f_i : S_i \to \text{Obj}(C)) \to (g : S_g \to \text{Obj}(C)) \). Define \( T : (f : \prod_{i \in I} S_i \to \text{Obj}(C)) \to (g : S_g \to \text{Obj}(C)) \) to be the function \( T(t, y) = \sum_{j \in I} t_j(y) \circ \text{Id}_j(y) = g(y) \).
The last equality also shows the uniqueness of $T$.

In particular we have the following.

**Corollary 1.2.** If $\emptyset \neq S_f \subset S_0$, then any object $f : S_f \to \Obj(C)$ is isomorphic to the direct sum of the objects $\{ f|_x : \{ x \} \to \Obj(C) \}_{x \in S}$.

Let us suppose now that the category $C$ is strict monoidal and that the set $S_0$ is infinite. In what follows we shall endow $C^{S_0}$ with a monoidal structure extending the one given in $C$. However, as we shall see, the structure that we define is not strict in general.

We start by defining the tensor product of objects and morphisms and a unit object. Next we define the associative constraint $A$, the left and right units $L$ and $R$, and finally we prove that they satisfy the required conditions.

First, we fix once and for all a bijection $\gamma : S_0 \times S_0 \to S_0$. Given two objects $f : S_f \to \Obj(C)$ and $g : S_g \to \Obj(C)$, define $f \otimes g$ by the following composite

$$f \otimes g : \gamma(S_f \times S_g) \xrightarrow{\gamma^{-1}} S_f \times S_g \xrightarrow{f \times g} \Obj(C) \times \Obj(C) \xrightarrow{\otimes} \Obj(C).$$

(6)

Choose any point $*$ on $S_0$ and define $I : \{ * \} \to \Obj(C)$ by $I(*) = I \in \Obj(C)$.

Now, for two morphisms $F : f \to f'$, $G : g \to g'$, and a point $(z, z') \in \gamma(S_f \times S_g) \times \gamma(S_{f'} \times S_{g'})$, define $F \otimes G : f \otimes g \to f' \otimes g'$ by

$$(F \otimes G)(z, z') := F(x_z, x'_{z'}) \otimes G(y_z, y'_{z'}) : f(x_z) \otimes g(y_z) \to f'(x'_{z'}) \otimes g'(y'_{z'}).$$

(7)

where $\gamma^{-1}(z) = (x_z, y_z) \in S_f \times S_g$ and $\gamma^{-1}(z') = (x'_{z'}, y'_{z'}) \in S_{f'} \times S_{g'}$ are the pairs such that $(f \otimes g)(z) = f(x_z) \otimes g(y_z)$ and $(f' \otimes g')(z') = f'(x'_{z'}) \otimes g'(y'_{z'})$.

It is clear that $\gamma(S^F_{x_z} \times S^G_{y_z}) \subset \gamma(S^F_{x_z} \times S^G_{y_z})$ is a finite set and that if $z' \in \gamma(S^F_{x_z} \times S^G_{y_z}) - \gamma(S^F_{x_z} \times S^G_{y_z})$, then $\gamma^{-1}(z') \notin S^F_{x_z} \times S^G_{y_z}$. Hence, either $x'_{z'} \notin S^F_{x_z}$ or $y'_{z'} \notin S^G_{y_z}$, and so $(F \otimes G)(z, z') = 0$ if $z' \notin \gamma(S^F_{x_z} \times S^G_{y_z})$.

Before we define the associative constraint $A$, we shall adopt the following notation. If, for example, $v \in \gamma(\gamma(S_f \times S_g) \times S_h)$, then we write

$$(\gamma^{-1} \times \id) \gamma^{-1}(v) = ((x_v, y_v), z_v) \in S_f \times S_g \times S_h.$$

Here, $\gamma(x_v, y_v)$ is the unique element in $\gamma(S_f \times S_g) \subset S_0$ such that $\gamma(\gamma(x_v, y_v), z_v) = v$. In other words, the inner parentheses will indicate the place from left to
right of the second $\gamma^{-1}$ in the composition $(\gamma^{-1} \times \text{id})\gamma^{-1}$. Analogously for $w \in \gamma(S_f \times \gamma(S_g \times S_h))$ we write

$$(\text{id} \times \gamma^{-1})(\gamma^{-1}(w)) = (x_w, (y_w, z_w)) \in S_f \times S_g \times S_h.$$  

When there is no risk of confusion we drop the inner parentheses and simply write $(\gamma^{-1} \times \text{id})\gamma^{-1}(v) = (x_v, y_v, z_v)$ and $(\text{id} \times \gamma^{-1})(\gamma^{-1}(w)) = (x_w, y_w, z_w)$. In the same way, if for example $v \in S_f \times S_g \times S_h$ then we write $$(\text{id} \times \gamma^{-1} \times \text{id})(\gamma^{-1}(v)) = ((x_v, (y_v, z_v)), t_v) \in S_f \times S_g \times S_h \times S_i,$$
or $(\text{id} \times \gamma^{-1} \times \text{id})(\gamma^{-1}(w)) = (x_v, y_v, z_v, t_v)$, etc.

With this notation, we have

$$(F \otimes G) \otimes H(v, w) = (F \otimes G)((x, y)_v, (x, y)_w) \otimes H(z_v, z_w) = F(x_v, x_w) \otimes G(y_v, y_w) \otimes H(z_v, z_w).$$  

(8)

Let us define $A_{f,g,h} : (f \otimes g) \otimes h \rightarrow f \otimes (g \otimes h)$ by

$A_{f,g,h}(v, w) = \delta^{w,v}_{x,y,z} \text{id}_{f(x_v, y_v, z_v) \otimes g(y_w, z_w)} : ((f \otimes g) \otimes h)(v) \otimes (f \otimes (g \otimes h))(w),$

(9)

where again, in order to shorten the notation, $\delta^{w,v}_{x,y,z}$ stands for $\delta_{x, y, z} \delta_{y, z, w} \delta_{z, w, v}$.

It is easy to see that the inverse of $A_{f,g,h}$ is given by

$A^{-1}_{f,g,h}(v, w) = \delta^{w,v}_{x,y,z} \text{id}_{f(x_v, y_v, z_v) \otimes g(y_w, z_w)} : (f \otimes (g \otimes h))(w) \otimes ((f \otimes g) \otimes h)(v)$.

(10)

Now we define the right unit $R_f : f \otimes I \rightarrow f$. For any object $f$, the object $f \otimes I$ is expressed by the composite

$$f \otimes I : \gamma(S_f \times \{\ast\}) \xrightarrow{\gamma^{-1}} S_f \times \{\ast\} \xrightarrow{f \times I} \text{Obj}(\mathcal{C}) \times \text{Obj}(\mathcal{C}) \xrightarrow{\otimes} \text{Obj}(\mathcal{C}).$$  

(11)

For $z \in \gamma(S_f \times \{\ast\})$, we write $\gamma^{-1}(z) = (x_z, \ast) \in S_f \times \ast$ and define $R_f : f \otimes I \rightarrow f$ by

$$R_f(z, x) = \delta_{x, x, z} \text{id}_{f(x_z)} : (f \otimes I)(z) = f(x_z) \rightarrow f(x)$$  

(12)

for $(z, x) \in \gamma(S_f \times \{\ast\}) \times S_f$. It is easy to see that $R_f$ is an isomorphism with inverse $R^{-1}_f : f \rightarrow f \otimes I$ given by the function

$$R^{-1}_f(x, z) = \delta_{x, x, z} \text{id}_{f(x)} : f(x) \rightarrow (f \otimes I)(z) = f(x_z).$$  

(13)

In the same way we define the left unit $L_f : I \otimes f \rightarrow f$, that is, if $z \in \gamma(\ast \times S_f)$, then we write $\gamma^{-1}(z) = (\ast, x_z)$ and define

$$L_f(z, x) = \delta_{x, x, z} \text{id}_{f(x_z)} : (I \otimes f)(z) = f(x_z) \rightarrow f(x)$$  

(14)

The inverse of $L_f$ is given by

$$L^{-1}_f(z, x) = \delta_{x, x, z} \text{id}_{f(x)} : f(x) \rightarrow (\text{id} \otimes f)(z) = f(x_z).$$  

(15)
Theorem 1.3. The category $C^S_0$ is a monoidal category with tensor product of objects and morphisms, associative constraint, and right and left units as we have just defined.

We divide the proof into four lemmas.

Lemma 1.4. If $F : f \to f'$, $F' : f' \to f''$, $G : g \to g'$, and $G' : g' \to g''$ are morphisms in $C^S_0$, then

(i) $(F' \otimes G') \circ (F \otimes G) = (F' \circ F) \otimes (G' \circ G)$ and

(ii) $\text{Id}_f \otimes \text{Id}_g = \text{Id}_{f \otimes g}$.

Proof. (i) For $z, z'' \in \gamma(S_f \times S_g)$ we have

$$((F' \otimes G') \circ (F \otimes G))(z, z'') = \sum_{z' \in \gamma(S_{f'} \times S_{g'})} (F' \otimes G')(z', z'') \circ (F \otimes G)(z, z')$$

$$= \sum_{z' \in \gamma(S_{f'} \times S_{g'})} (F'(x'_z, x'_{z''}) \circ F(x_z, x'_z)) \otimes (G'(y'_z, y'_{z''}) \circ G(y_z, y'_z))$$

$$= \left( \sum_{x' \in S_{f'}} F'(x', x''_{z''}) \circ F(x_z, x') \right) \otimes \left( \sum_{y' \in S_{g'}} G'(y', y''_{z''}) \circ G(y_z, y') \right)$$

$$= (F' \circ F)(x_z, x''_{z''}) \otimes (G' \circ G)(y_z, y''_{z''})$$

$$= ((F' \circ F) \otimes (G' \circ G))(z, z'')$$

The third equality follows from the fact that $\gamma$ establishes a bijection between $S_{f'} \times S_{g'}$ and $\gamma(S_{f'} \times S_{g'})$.

(ii) For $z, z' \in \gamma(S_f \times S_g)$ we have

$$(\text{Id}_f \otimes \text{Id}_g)(z, z') = \text{Id}_f(x_z, x'_{z'}) \otimes \text{Id}_g(y_z, y_{z'})$$

$$= \delta_{x_z, x'} \text{Id}_f(x_z) \otimes \delta_{y_z, y_{z'}} \text{Id}_g(y_z)$$

$$= \delta_{x_z, x'} \text{Id}_f(x_z) \otimes \text{Id}_g(y_z)$$

$$= \delta_{x_z, x'} \text{Id}_f(g(y_z))$$

$$= \text{Id}_{f \otimes g}(z, z')$$

Lemma 1.5. The associative constraint $A$ defined above is a natural isomorphism that satisfies the Pentagonal Axiom.
Proof. We already saw that $A$ is an isomorphism. To show that it is natural, we have, on the one hand,

\[
(F \otimes (G \otimes H)) \circ A_{f,g,h}(v, w') = \sum_{w \in \gamma(S_f \times \gamma(S_y \times S_h))} (F \otimes (G \otimes H))(w, w') \circ A_{f,g,h}(v, w)
\]

\[
= \sum_{w \in \gamma(S_f \times \gamma(S_y \times S_h))} (F(x_w, x_{w'}) \otimes G(y_v, y_{w'})) \otimes H(z_w, z_{w'}) \circ \delta_{x,y,z}^{i,j,k} \circ \gamma_{f,g,h}(v) \circ \gamma_{g,h}(y) \circ \gamma_{h}(z)
\]

\[
= F(x_v, x_{w'}) \otimes G(y_v, y_{w'}) \otimes H(z_v, z_{w'}).
\]

On the other hand, we have

\[
(A_{f',g',h'} \circ ((F \otimes G) \otimes H))(v, w') = \sum_{w' \in \gamma(S_{f'} \times \gamma(S_y \times S_{h'}))} A_{f',g',h'}(v', w') \circ ((F \otimes G) \otimes H)(v, v')
\]

\[
= \sum_{w' \in \gamma(S_{f'} \times \gamma(S_y \times S_{h'}))} \delta_{x,y,z}^{i,j,k} \circ \gamma_{f',g',h'}(v') \circ \gamma_{g',h'}(y) \circ \gamma_{h'}(z)
\]

\[
= F(x_v, x_{w'}) \otimes G(y_v, y_{w'}) \otimes H(z_v, z_{w'}).
\]

Therefore, $A_{f',g',h'} \circ ((F \otimes G) \otimes H) = (F \otimes (G \otimes H)) \circ A_{f,g,h}$, so $A$ is natural.

Let us prove now that $A$ satisfies the Pentagonal Axiom. Set $M(s, w) = ((id_f \otimes A_{g,h,i}) \circ A_{f,g,h,i} \circ (A_{f,g,h} \otimes id_i))(s, w)$. For $s \in \gamma(\gamma(S_f \times S_y \times S_h) \times S_i)$ and $w \in \gamma(\gamma(S_f \times \gamma(S_y \times \gamma(S_h \times S_i))))$, we have

\[
M(s, w) = \sum_{u \in S_f \otimes ((y \otimes h) \otimes i)} (id_f \otimes A_{g,h,i})(v, w) \circ A_{f,g,h,i}(u, v) \circ (A_{f,g,h} \otimes id_i)(s, u)
\]

\[
= \sum_{u \in S_f \otimes ((y \otimes h) \otimes i)} (\delta_{x, y, z}^{w, u} \circ \delta_{y, z, t}^{w, v} \circ \delta_{y, z, t}^{w, v}) \circ \delta_{x, y, z, t}^{w, v} \circ \gamma_{f,g,h}(v) \circ \gamma_{g,h}(y) \circ \gamma_{h}(z)
\]

\[
= \delta_{x,y,z,t}^{w} \circ \text{id}_{f(x_u) \otimes g(y_u) \otimes h(z_u)} \circ \delta_{x,y,z,t}^{w} \circ \gamma_{f,g,h}(v) \circ \gamma_{g,h}(y) \circ \gamma_{h}(z)
\]

\[
= \delta_{x,y,z,t}^{w} \circ \text{id}_{f(x_u) \otimes g(y_u) \otimes h(z_u)} \circ \delta_{x,y,z,t}^{w} \circ \gamma_{f,g,h}(v) \circ \gamma_{g,h}(y) \circ \gamma_{h}(z).
\]

(20)
Set \( N(s, w) = (A_{f,g,h,i} \circ A_{f,g,h,i})(s, w) \). Then,

\[
N(s, w) = \sum_{r \in S(f,g)} (A_{f,g,h,i})(r, w) \circ (A_{f,g,h,i})(s, r) \\
= \sum_{r \in S(f,g)} \delta^{x,w}_{x,y,z} \text{id}_r (x_t) \otimes (y_t) \circ \delta^{x,r}_{x,y,z} \text{id}_s (x_s) \otimes (y_s) \\
= \delta^{x,w}_{x,y,z} \text{id}_r (x_t) \otimes (y_t) \otimes \text{id}_s (x_s) \otimes (y_s). \tag{21}
\]

Thus \( M(s, w) = N(s, w) \) and so, \( A \) satisfies the Pentagonal Axiom.

**Lemma 1.6.** The right unit \( R \) and the left unit \( L \) are natural isomorphisms.

**Proof.** We already saw that \( R_f \) is an isomorphism. For \( z \in \gamma(S_f \times *) \) and \( x' \in S_f \), we have

\[
(F \circ R_f)(z, x') = \sum_{x \in S_f} F(x, x') \circ R_f(z, x) \\
= F(x, x') \circ \delta_{x,x'} \text{id}_f(x_i) \\
= F(x_z, x'). \tag{22}
\]

On the other hand

\[
(R_f \circ (F \otimes \text{id}_l))(z, x') = \sum_{z' \in \gamma(S_f \times *)} R_f(z', x') \circ (F \otimes \text{id}_l)(z, z') \\
= \delta_{x,x'} \text{id}_f(x_z) \circ (F(x_z, x_z') \otimes \text{id}_l(*, *)) \\
= \delta_{x,x'} \text{id}_f(x_z) \circ F(x_z, x_z') \\
= F(x_z, x'). \tag{23}
\]

The proof for \( L \) is analogous.

**Lemma 1.7.** The morphisms \( A, R \) and \( L \) satisfy the Triangular Axiom.

**Proof.** Set \( P(v, w) = ((\text{id}_f \otimes L_g) \circ A_{f,i,g})(v, w) \). Then

\[
P(v, w) = \sum_{u \in S(f,g)} (\text{id}_f \otimes L_g)(u, w) \circ A_{f,i,g}(v, u) \\
= (\delta_{x,v,x,w} \text{id}_f(x_z) \otimes \delta_{y,v,y,w} \text{id}_g(y_z)) \circ \delta^{x,w}_{x,y,z} \text{id}_f(x_z) \otimes \text{id}_g(y_z) \\
= \delta_{x,v,x,w} \text{id}_f(x_z) \otimes \delta_{y,v,y,w} \text{id}_g(y_z) \\
= (R_f \otimes \text{id}_g)(v, w). \tag{24}
\]

So \( (\text{id}_f \otimes L_g) \circ A_{f,i,g} = R_f \otimes \text{id}_g \).

These four lemmas finish the proof of \( 1 \).\( \square \)
**Proposition 1.8.** The category $C^S_0$ has a full subcategory, which is tensor equivalent to $C$.

**Proof.** Recall that a tensor functor is a triple $(F, \varphi_0, \varphi_2)$, where $F$ is a functor, $\varphi_0$ is an isomorphism from $I$ to $F(I)$, and $\varphi_2(U, V) : F(U \otimes F(V) \rightarrow F(U \otimes V)$ is a family of natural isomorphisms compatible with the associative constraint and the left and right units (see [1] p.287). Define a functor $J : C \rightarrow C^S_0$, by choosing for any object $V$ in $C$ any point $x_V \in S_0$ and a function $f_V : \{x_V\} \rightarrow \text{Obj}(C)$, given by $f_V(x_V) = V$. Then we define $J(V) = f_V$. To any morphism $\alpha : V \rightarrow W$ we assign the function $F_\alpha(x_V, x_W) = \alpha : f_V(x_V) = V \rightarrow f_W(x_W) = W$ and then define $J(\alpha) = F_\alpha$. For the unit object $I$ of $C$ we choose the fixed point $\ast$ as before, so that $J(I) = I \in \text{Obj}(C^S_0)$. For $U, V$ objects of $C$, define $\varphi_2(U, V) : J(U) \otimes J(V) = f_U \otimes f_V \rightarrow J(U \otimes V) = f_{U \otimes V}$, as follows. If $\gamma^{-1}(\{x_V\} \times \{x_V\}) = \{x'_U, V\}$, then $(f_U \otimes f_V)(x'_U, V) = U \otimes V$ and $f_{U \otimes V}(x'_U, V) = U \otimes V$, then take $\varphi_2(U, V)(x'_U, V) = \text{id}_{U \otimes V}$. The morphisms $\varphi_0$ and $\varphi_2$ are equivalent to $C$, it is clear that $C$ is invertible with inverse given by $C^{-1}(w, v) = \delta^{w, v}_{x, y} c^{-1}_{f(x_w), g(y_v)}$.

**Proposition 1.9.** The family $C$ of isomorphisms $C_{f, g}$ is a braiding in the category $C^S_0$.

**Proof.** We have to prove that $C$ is natural and satisfies the Hexagonal Axiom. For $F : f \rightarrow f'$ and $G : g \rightarrow g'$ we have, on the one hand

\[
(G \otimes F) \circ C_{f, g}(v, w) = \sum_{w' \in \gamma(S_g \otimes S_f)} (G \otimes F)(w, w') \circ C_{f, g}(v, w)
\]

\[
= \sum_{w' \in \gamma(S_g \otimes S_f)} (G(y_{w'}, y_{w'}) \otimes F(x_{w'}, x_{w'})) \circ \delta^{w, w'}_{x, y} c_{f(x_w), g(y_v)}
\]

\[
= (G(y_v, y_{w'}) \otimes F(x_v, x_{w'})) \circ c_{f(x_v), g(y_v)}.
\]

(26)
On the other hand,

\[
C_{f',g'} (F \otimes G)(v, w') = \sum_{v' \in \gamma(S_f \times S_{f'})} C_{f',g'}(v', w') \circ (F \otimes G)(v, v')
\]
\[
= \sum_{v' \in \gamma(S_f \times S_{f'})} \delta_{v',w'} \delta_{f',g'}(y_{v'}) \circ F(x_v, x_{v'}) \otimes G(y_v, y_{v'})
\]
\[
= c_{f'(x_{v'}),g'(y_{v'})} \circ (F(x_v, x_{v'}) \otimes G(y_v, y_{v'})).
\] (27)

Both sums are equal since \( c \) is a braiding in \( \mathcal{C} \) and therefore it is natural. Thus \( \mathcal{C} \) is natural. We now show the commutativity of one of the diagrams of the Hexagonal Axiom. Put \( M(w, w') = (A_{g,h,f} \circ C_{f,g,h} \circ A_{g,h})(w, w') \). Then

\[
M(w, w') = \sum_{u \in S_f \otimes (g \otimes h)} A_{f,g,h}(u, w') \circ C_{f,g,h}(v, u) \circ A_{f,g,h}(w, v)
\]
\[
= \sum_{u \in S_f \otimes (g \otimes h)} \delta_{u,w'} \delta_{f,g,h}(y_u) \circ \delta_{f,g,h}(x_u) \circ \delta_{f,g,h}(z_u)
\] (28)

Set \( N(w, w') = ((\text{Id}_g \otimes C_{f,h}) \circ A_{g,f,h} \circ (C_{f,g} \otimes \text{Id}_h))(w, w') \). Then

\[
N(w, w') = \sum_{u \in S_g \otimes (f \otimes h)} (\text{Id}_g \otimes C_{f,h})(u, w') \circ A_{g,f,h}(v, u) \circ (C_{f,g} \otimes \text{Id}_h)(w, v)
\]
\[
= \sum_{u \in S_g \otimes (f \otimes h)} \delta_{u,w'} \delta_{g,h}(y_u) \circ \delta_{f,g,h}(x_u) \circ \delta_{f,g,h}(z_u)
\]
\[
= \delta_{u,w'} \delta_{g,h}(y_u) \circ \delta_{f,g,h}(x_u) \circ \delta_{f,g,h}(z_u).
\] (29)

Again, since \( c \) is a strict braiding in \( \mathcal{C} \), we have the equality \( M(w, w') = N(w, w') \). The commutativity of the other hexagon is proved analogously. \( \square \)

In the same way, if the category \( \mathcal{C} \) has a twist, then we can easily prove the following assertion.

**Proposition 1.10.** Let \( \theta \) be a twist for the the category \( \mathcal{C} \). Then the category \( \mathcal{C}^{\otimes} \) has a twist \( \Theta_f : f \rightarrow f \) given by

\[
\Theta_f(x, y) = \delta_{x,y} \theta_f(x) : f(x) \rightarrow f(y)
\] (30)

for any object \( f \) in \( \mathcal{C}^{\otimes} \). \( \square \)
However, it is not possible to extend a duality from \( \mathcal{C} \) to \( \mathcal{C}^{S_0} \). Although we have for any \( f : S_f \to \text{Obj}(\mathcal{C}) \) a canonical candidate for \( f^* : S_f \to \text{Obj}(\mathcal{C}) \), namely the function \( f^* \) defined by \( f^*(x) = (f(x))^* \) as well as a canonical candidate for the evaluation \( D_f : f^* \otimes f \to I \), given by \( D_f(v, \{s\}) = \delta_{x^*_v, x_v} d_f(x_v) : f^*(x^*_v) \otimes f(x_v) \to I(\{s\}) = I \), where \( \gamma^{-1}(v) = (x^*_v, x_v) \in S_f \times S_f \), this is not the case for the coevaluation. Indeed, the canonical extension \( B_f : I \to f \otimes f^* \) given by \( B_f(s, v) = \delta_{x^*_v, x_v} b_f(x_v) : I \to f(x_v) \otimes f^*(x_v) \) is not a morphism in \( \mathcal{C}^{S_0} \) if \( S_f \) is infinite, since condition (ii) of page 2 does not hold.

Nevertheless, if we consider the full subcategory \( \mathcal{C}^{S_0}_2 \) which as objects has functions \( f \) with finite domain \( S_f \), then it is possible to extend the duality according to the given formulas. It is easy to see that the inclusion functor \( J : \mathcal{C} \to \mathcal{C}^{S_0}_2 \) factors through \( \mathcal{C}^{S_0}_3 \), i.e.,

\[
\begin{array}{ccc}
\mathcal{C} & \xrightarrow{J} & \mathcal{C}^{S_0}_2 \\
\mathcal{C}^{S_0}_2 & \xrightarrow{\mathcal{C}^{S_0}} & \mathcal{C}^{S_0}
\end{array}
\]

The following assertion is also easy to prove.

**Proposition 1.11.** If the category \( \mathcal{C} \) is a ribbon category, then the extended structure in \( \mathcal{C}^{S_0}_2 \) is pivotal braided (but nonstrict in general, so it is not ribbon).

**Remark 1.12.** In order to simplify the next computations, we shall adopt the following notation. Let \( \mathcal{A} \) be the set of isomorphisms of \( \mathcal{C}^{S_0} \) generated by the set \( \{\text{Id}_{\chi}, A^\pm_{\kappa, \lambda, \mu}, R_\zeta, L_\varsigma\} \) under tensor products and compositions, where \( \chi, \kappa, \lambda, \mu, \zeta, \) and \( \varsigma \) are any objects in \( \mathcal{C}^{S_0} \). In other words, \( \mathcal{A} \) is the set of isomorphisms that relate different objects by associativity and units. If \( F \) and \( G \) are morphisms in \( \mathcal{C}^{S_0} \), we shall write \( F \doteq G \) if \( G = X \circ F \circ Y \), where \( X \) and \( Y \) are elements of \( \mathcal{A} \). For example, \( F \doteq G \) if the following diagram commutes.

\[
\begin{array}{c}
((f_1 \otimes f_2) \otimes f_3) \otimes f_5 \xrightarrow{F} (g_1 \otimes g_2) \otimes g_3 \\
A_{f_1 \otimes f_2, f_3, f_5} \otimes \text{id}_{f_5} \\
((f_1 \otimes f_2) \otimes (f_3 \otimes f_4)) \otimes f_5 \\
A_{f_1, f_2, f_3, f_4, f_5} \\
(f_1 \otimes f_2) \otimes ((f_3 \otimes f_4) \otimes f_5) \\
\text{id}_{f_1} \otimes f_2 \otimes A_{f_3, f_4, f_5}^{-1} \\
(f_1 \otimes f_2) \otimes (f_3 \otimes (f_4 \otimes f_5)) \xrightarrow{G} g_1 \otimes (g_2 \otimes g_3)
\end{array}
\]

The relation \( \doteq \) is an equivalence relation in the set of morphisms of \( \mathcal{C}^{S_0} \) which is compatible with composition and tensor product in the sense that if \( F \doteq G \) and \( F' \doteq G' \) then \( F \circ F' \doteq G' \circ G \), if the compositions are defined, and \( F \otimes F' \doteq G \otimes G' \). Indeed, for the composition, suppose \( A \circ F \circ B = G \) and
that $C \circ F' \circ D = G'$, for elements $A$, $B$, $C$, and $D$ in $\mathcal{A}$. Then $G' \circ G = C \circ F' \circ D \circ A \circ F \circ B$. The morphism $D \circ A$ is an endomorphism of the domain $s(F')$ of $F'$ which is equal to the codomain $t(F)$ of $F$ and is an element of $\mathcal{A}$. Mac Lane’s coherence theorem states that this element has to be the identity morphism $\text{Id}_{s(F')}$. Hence $G' \circ G = C \circ F' \circ F \circ B$. The tensor part follows from the identity $(A \circ F \circ B) \otimes (C \circ F' \circ D) = (A \otimes C) \circ (F \otimes F') \circ (B \otimes D)$. In what follows we shall use this notation without further comments.

2 Bialgebras in $\mathcal{C}^{S_0}$

Let $\mathcal{V}$ be a monoidal category. We say that an object $A$ of $\mathcal{V}$ is an algebra in $\mathcal{V}$, if there exist morphisms $\mu : A \otimes A \to A$ and $\eta : I \to A$ such that

$$\mu(\mu \otimes \text{id}_A) = \mu(\text{id}_A \otimes \mu), \quad (32)$$

$$\mu(\eta \otimes \text{id}_A) = \text{id}_A = \mu(\text{id}_A \otimes \eta). \quad (33)$$

Dually, we say that $C$ is a coalgebra in $\mathcal{V}$, if there exist morphisms $\Delta : C \to C \otimes C$ and $\varepsilon : C \to I$ such that

$$\Delta(\Delta \otimes \text{id}_C) \Delta = (\text{id}_C \otimes \Delta) \Delta, \quad (34)$$

$$\varepsilon(\varepsilon \otimes \text{id}_C) \Delta = \text{id}_C \Delta = (\text{id}_C \otimes \varepsilon) \Delta. \quad (35)$$

If $H$ is an algebra, then the product in $H \otimes H$ is defined by the following composite

$$\tilde{\mu} : (H \otimes H) \otimes (H \otimes H) \xrightarrow{A_{H, H, H}^{-1}} ((H \otimes H) \otimes H) \otimes H \xrightarrow{A_{H, H, H \otimes \text{id}_H}} (H \otimes H) \otimes H \otimes (H \otimes H) \xrightarrow{\mu \otimes \mu} H \otimes H.$$ (36)

We say that $H$ is a bialgebra in $\mathcal{V}$, if $\tilde{\mu}(\Delta \otimes \Delta) = \Delta \mu$ and $\varepsilon \mu = \varepsilon \otimes \varepsilon$.

If $A$ is an algebra, an object $V$ is an $A$-module, if there exists a morphism $T : A \otimes V \to V$, such that $T(\mu \otimes \text{id}_V) = T(\text{id}_A \otimes T)$ and $T(\eta \otimes \text{id}_V) = \text{id}_V$.

Note that if the category is strict monoidal, the latter are the concepts of algebra, coalgebra, bialgebra and module in strict braided monoidal categories.

We are going to find bialgebras in $\mathcal{C}^{S_0}$, when $\mathcal{C}$ is a braided strict monoidal category with left duality.

Let $h : S_h \to \text{Obj}(\mathcal{C})$ be an injective function such that $h(S_h) \subset \text{Obj}(\mathcal{C})$ is closed under $\otimes$, that is, for any pair $(x, y) \in S_h \times S_h$, there exists a unique
$z \in S_h$ such that $h(x) \otimes h(y) = h(z)$ and suppose $I \in h(S_h)$. For example, we can take a set $S_0$ with the same cardinality as $\text{Obj}(C)$ and $h : S_0 \rightarrow \text{Obj}(C)$ to be any bijection, if $\text{Obj}(C)$ is an infinite set.

Set $\Delta_h = \{(x, x) \mid x \in S_h\} \subset S_h \times S_h$ and let $\overline{h}$ be the object defined by the composite

$$\overline{h} : \gamma(\Delta_h) \xrightarrow{\gamma^{-1}} \Delta_h h^* \otimes h \text{Obj}(C) \times \text{Obj}(C) \xrightarrow{\otimes} \text{Obj}(C)$$

(37)

where $h^*(x) := (h(x))^*$. That is, $\overline{h}$ is defined by the relation $\overline{h}(\gamma(x, x)) = h^*(x) \otimes h(x)$, for $\gamma(x, x) \in S_h = \gamma(\Delta_h)$.

The main theorem in this section is the following.

**Theorem 2.1.** The object $\overline{h}$ is a bialgebra in $C_{S_0}$ and the objects of $C$, considered as a subcategory of $C_{S_0}$, are $\overline{h}$-modules.

To prove it, we shall establish two previous lemmas. Let $\chi : S_h \times S_h \rightarrow S_h$ be the function defined by the relation $h(\chi(x, y)) = h(x) \otimes h(y)$.

**Lemma 2.2.** The function $\chi$ satisfies $\chi(\chi(x, y), z) = \chi(x, \chi(y, z))$.

**Proof.**

$$h(\chi(x, y), z)) = h(\chi(x, y)) \otimes h(z) = h(x) \otimes h(y) \otimes h(z) = h(x) \otimes h(\chi(y, z)) = h(x, \chi(y, z)).$$

Thus $\chi(\chi(x, y), z) = \chi(x, \chi(y, z)).$ \hfill \qed

We now consider, in a more general situation, a monoidal category with left duality $V$. In the following lemma we use letters $x, y, z, \ldots$ to denote objects of $V$. Let $x, y$ be objects of $V$. Recall that there exists an isomorphism $\gamma_{x, y} : y^* \otimes x^* \rightarrow (x \otimes y)^*$ given by

$$\gamma_{x, y} = (\text{id}_y \otimes \text{id}_{(x \otimes y)^*})(\text{id}_{y^*} \otimes \text{id}_x \otimes \text{id}_{(x \otimes y)^*})(\text{id}_{y^*} \otimes x^* \otimes \text{id}_{x \otimes y}).$$

(38)

Now define the isomorphism $\Gamma_{x, y} : y^* \otimes y \otimes x^* \otimes x \rightarrow (x \otimes y)^* \otimes (x \otimes y)$ by the composite

$$\Gamma_{x, y} : y^* \otimes y \otimes x^* \otimes x \xrightarrow{\text{id}_{x^*} \otimes \gamma_{x, y} \otimes \text{id}_x} y^* \otimes x^* \otimes y \otimes x \xrightarrow{\gamma_{x, y}} (x \otimes y)^* \otimes (x \otimes y).$$

**Lemma 2.3.** The isomorphisms $\Gamma_{x, y}$ satisfy the relation

$$\Gamma_{x, y \otimes z}(\Gamma_{y, z} \otimes \text{id}_{x^* \otimes x}) = \Gamma_{x \otimes y, z}(\text{id}_{y^* \otimes z} \otimes \Gamma_{x, y}).$$

That is, if $x, y$ and $z$ are objects of $V$ then the following diagram commutes

$$\begin{array}{ccc}
z^* \otimes z \otimes y^* \otimes y \otimes x^* \otimes x & \xrightarrow{\Gamma_{x, y \otimes z} \otimes \text{id}_{x^* \otimes x}} & (y \otimes x)^* \otimes (y \otimes x) \otimes x^* \otimes x \\
\text{id}_{y^* \otimes z} \otimes \Gamma_{x, y} & & \Gamma_{x, y \otimes z} \\
z^* \otimes z \otimes (x \otimes y)^* \otimes (x \otimes y) & \xrightarrow{\Gamma_{x \otimes y, z}} & (x \otimes y \otimes z)^* \otimes (x \otimes y \otimes z).
\end{array}$$
Proof. We prove it by using graphical calculus. In Figure 1 the morphism \( \Gamma_{x,y} \) is represented. Figure 2 (at the end of the paper) proves the Lemma. The upper left and bottom right diagrams represent the morphisms \( \Gamma_{x,y} \otimes z (\Gamma_{y,z} \otimes \text{id}_x \otimes \text{id}_x) \) and \( \Gamma_{x \otimes y,z} (\text{id}_z \otimes \text{id}_z \otimes \Gamma_{x,y}) \), respectively.

\[
\Gamma_{x,y} = \begin{array}{c}
\text{id}_x \otimes y \\
\end{array}
\]

Figure 1: The morphism \( \Gamma_{x,y} \)

\[\text{(39)}\]

**Proof of Theorem 2.1** Define \( \mu : \overline{h} \otimes \overline{h} \rightarrow \overline{h} \) by

\[
\mu(v, \gamma(z, z)) : (\overline{h} \otimes \overline{h})(v) = h^*(x_v) \otimes h(x_v) \otimes h^*(y_v) \otimes h(y_v) \\
\delta_{x,v,x_v} \Gamma_{h(y_v), h(x_v)} \rightarrow h^*(z) \otimes h(z)
\]

Since there is a unique \( x_0 \in S_h \) such that \( h(x_0) = I \in \text{Obj}(C) \), we can define \( \eta : I \rightarrow \overline{h} \) by

\[
\eta(\ast, \gamma(y, y)) = \delta_{x_0, y} \text{id}_I : I = h^*(x_0) \otimes h(x_0) \rightarrow h^*(y) \otimes h(y).
\]

We have to prove now that \( \mu(\mu \otimes \text{id}_\overline{h}) = \mu(\text{id}_\overline{h} \otimes \mu) \) and \( \mu(\eta \otimes \text{id}_\overline{h}) = \text{id}_\overline{h} = \mu(\text{id}_\overline{h} \otimes \eta) \).

Set \( S = \mu(\mu \otimes \text{id}_\overline{h})(w, \gamma(t, t)) \) and \( R = \mu(\text{id}_\overline{h} \otimes \mu)(w', \gamma(t, t)) \). We have on the
one hand

\[ S = \sum_{v \in S_{\mu,v}} \mu(v, \gamma(t,t)) \circ (\mu \otimes \text{Id}_{\tilde{T}})(w, v) \]
\[ = \sum_{v \in S_{\mu,v}} \delta_t,\chi(y_v, x_v) \Gamma h(y_v), h(x_v) \circ (\delta_{x_v,\Gamma h(y_v), h(x_v)} \Gamma h(y_v), h(x_v) \otimes \delta_{z_v, y_v} \text{id}_{h^*(z_v) \otimes h(z_v)}) \]
\[ = \delta_t,\chi(z_v, \chi(y_v, x_v)) \Gamma h(z_v), h(x_v) \circ (\Gamma h(y_v), h(x_v) \otimes \text{id}_{h^*(z_v) \otimes h(z_v)}) \]
\[ = \delta_t,\chi(z_v, \chi(y_v, x_v)) \Gamma h(z_v), h(y_v) \circ (\Gamma h(y_v), h(x_v) \otimes \text{id}_{h^*(z_v) \otimes h(z_v)}) \]  

(40)

On the other hand we have

\[ R = \sum_{v \in S_{\mu,v}} \mu(v, \gamma(t,t)) \circ (\text{Id}_{\tilde{T}} \otimes \mu)(w', v) \]
\[ = \sum_{v \in S_{\mu,v}} \delta_t,\chi(y_v, x_v) \Gamma h(y_v), h(x_v) \circ (\delta_{x_v,\Gamma h(y_v), h(x_v)} \Gamma h(y_v), h(x_v) \otimes \delta_{z_v, y_v} \text{id}_{h^*(z_v) \otimes h(z_v)}) \]
\[ = \delta_t,\chi(z_v, \chi(y_v, x_v)) \Gamma h(z_v), h(x_v) \circ (\Gamma h(y_v), h(x_v) \otimes \text{id}_{h^*(z_v) \otimes h(z_v)}) \]
\[ = \delta_t,\chi(z_v, \chi(y_v, x_v)) \Gamma h(z_v), h(y_v) \circ (\Gamma h(y_v), h(x_v) \otimes \text{id}_{h^*(z_v) \otimes h(z_v)}) \]  

(41)

According to Lemma 2.22 we have \( \chi(\chi(z_w', y_w'), x_w') = \chi(z_w', \chi(y_w', x_w')). \)

From this and Lemma 2.3 it is easy to see that \( R \circ A_{\mu,\tilde{H},\tilde{H}} = S \) so \( R \equiv S. \)

We shall prove now that \( \mu(\eta \otimes \text{Id}_{\tilde{T}}) \equiv \text{Id}_{\tilde{T}}. \) Set \( J = \mu(\eta \otimes \text{Id}_{\tilde{T}})(u, \gamma(z, z)). \)

From \( h(\chi(x_u, x_0)) = h(x_u) \otimes h(x_0) = h(x_u) \otimes \text{Id} = h(x_u) \) we deduce that \( \chi(x_u, x_0) = x_u \) and since \( \Gamma_{a,1} = \text{id}_{h^*(a) \otimes h(a)} \) for any object \( a \) of \( C, \) we have

\[ J = \sum_{v \in S_{\mu,v}} \mu(v, \gamma(z, z)) \circ (\eta \otimes \text{Id}_{\tilde{T}})(u, v) \]
\[ = \sum_{v \in S_{\mu,v}} \delta_z,\chi(y_v, x_v) \Gamma h(y_v), h(x_v) \circ (\eta(\gamma(z, z), \chi(y_v, x_v))) \otimes \text{Id}_{\tilde{T}}(\gamma(z, z), \chi(y_v, x_v))) \]
\[ = \sum_{v \in S_{\mu,v}} \delta_z,\chi(y_v, x_v) \Gamma h(y_v), h(x_v) \circ (\delta_{x_v,\eta} \text{id}_{\tilde{T}} \otimes \delta_{x_v,y_v} \text{id}_{h^*(x_v) \otimes h(x_v)}) \]
\[ = \delta_z,\chi(x_u, x_0) \Gamma h(x_u), h(x_0) \]
\[ = \delta_{z,u} \Gamma h(x_u), h(x_0) \]
\[ = \delta_z,\text{id}_{h^*(x_v) \otimes h(x_v)} \]
\[ = \text{Id}_{\tilde{T}}(u, \gamma(z, z)) \]  

(42)

The relation \( \mu(\text{Id}_{\tilde{T}} \otimes \eta) \equiv \text{Id}_{\tilde{T}} \) is proved in a similar way.

We have thus shown that \((\tilde{H}, \mu, \eta)\) is an algebra in \( C^{S_0}. \) Define now \( \Delta : \tilde{H} \rightarrow \tilde{H} \otimes \tilde{H} \) to be the function

\( \Delta(\gamma(x, x), v) : h^*(x) \otimes h(x) \rightarrow h^*(y_v) \otimes h(y_v) \otimes h^*(z_v) \otimes h(z_v) \)

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given by the following composite
\[
\delta_{x,y} \delta_{x,z} \alpha_{h(x)} \otimes b_{h(x)} \otimes \alpha_{h(x)} : h^*(x) \otimes h(x) \rightarrow h^*(y) \otimes h(y) \otimes h^*(z) \otimes h(z)
\]
and define \( \varepsilon : \mathbb{H} \rightarrow I \) as the function given by
\[
\varepsilon(\gamma(x,x), \ast) = d_{h(x)} : h^*(x) \otimes h(x) \rightarrow I.
\]
We are going to prove that \((\text{Id}_\mathbb{H} \otimes \Delta) \Delta = \Delta \otimes \text{Id}_\mathbb{H} \Delta\) and \((\varepsilon \otimes \text{Id}_\mathbb{H}) \Delta = \text{Id}_\mathbb{H} = (\text{Id}_\mathbb{H} \otimes \varepsilon)\). Set \( L = (\text{Id}_\mathbb{H} \otimes \Delta)(\gamma(t,t),w) \). Then
\[
L = \sum_{v \in S_{\mathbb{H}}} (\text{Id}_\mathbb{H} \otimes \Delta)(v, w') \circ \Delta(\gamma(t,t), v)
\]
\[
= \sum_{v \in S_{\mathbb{H}}} \delta_{x,v} \delta_{x,w} \alpha_{h(x)} \otimes b_{h(x)} \otimes \alpha_{h(x)} \otimes \delta_{y,z} \delta_{y,w} \alpha_{h(y)} \otimes b_{h(y)} \otimes \alpha_{h(y)} \circ
\]
\[
(\delta_{t,x} \delta_{t,w} \alpha_{h(t)} \otimes b_{h(t)} \otimes \alpha_{h(t)}
\]
\[
= \delta_{t,x} \delta_{t,w} \alpha_{h(t)} \otimes b_{h(t)} \otimes \alpha_{h(t)} \circ (\alpha_{h(t)} \otimes b_{h(t)} \otimes \alpha_{h(t)})
\]
(43)

Set \( R = (\Delta \otimes \text{Id}_\mathbb{H} \Delta)(\gamma(t,t),w) \). Then
\[
R = \sum_{v \in S_{\mathbb{H}}} (\Delta \otimes \text{Id}_\mathbb{H} \Delta)(v, w) \circ \Delta(\gamma(t,t), v)
\]
\[
= \sum_{v \in S_{\mathbb{H}}} \delta_{x,v} \delta_{x,w} \alpha_{h(x)} \otimes b_{h(x)} \otimes \alpha_{h(x)} \otimes \delta_{y,z} \delta_{y,w} \alpha_{h(y)} \otimes b_{h(y)} \otimes \alpha_{h(y)} \circ
\]
\[
(\delta_{t,x} \delta_{t,w} \alpha_{h(t)} \otimes b_{h(t)} \otimes \alpha_{h(t)}
\]
\[
= \delta_{t,x} \delta_{t,w} \alpha_{h(t)} \otimes b_{h(t)} \otimes \alpha_{h(t)} \circ (\alpha_{h(t)} \otimes b_{h(t)} \otimes \alpha_{h(t)})
\]
(44)

Taking \( x = h(t) \), Figure \( \Box \) (at the end of the paper) shows that \( R \) and \( L \) are equal up to associativity, that is \( A_{\mathbb{H},\mathbb{H},\mathbb{H}}(w, w') \circ R = L \). Thus \( L = R \).

Next we prove \((\varepsilon \otimes \text{Id}_\mathbb{H}) \Delta = \text{Id}_\mathbb{H} \). Set \( J = (\varepsilon \otimes \text{Id}_\mathbb{H}) \Delta(\gamma(x,x), \gamma(y,y)) \). Then
\[
J = \sum_{x \in S_{\mathbb{H}}} (\varepsilon \otimes \text{Id}_\mathbb{H})(v, \gamma(y,y)) \circ \Delta(\gamma(x,x), v)
\]
\[
= \sum_{x \in S_{\mathbb{H}}} \varepsilon(\gamma(x,x), \ast) \otimes \text{Id}_\mathbb{H}(\gamma(y,y), \gamma(y,y)) \circ
\]
\[
(\delta_{x,x} \delta_{x,y} \alpha_{h(x)} \otimes b_{h(x)} \otimes \alpha_{h(x)}
\]
\[
= \delta_{x,x} \delta_{x,y} \alpha_{h(x)} \otimes b_{h(x)} \otimes \alpha_{h(x)} \circ (\alpha_{h(x)} \otimes b_{h(x)} \otimes \alpha_{h(x)})
\]
(45)

From the definition of left duality we get \((d_{h(x)} \otimes \alpha_{h(x)})(\alpha_{h(x)} \otimes b_{h(x)}) = \alpha_{h(x)} \), so \( J = \delta_{x,y} \alpha_{h(x)} \otimes b_{h(x)} = \text{Id}_\mathbb{H}(\gamma(x,x), \gamma(y,y)) \).
The relation $\text{Id}_\overline{\Pi} = (\text{Id}_\overline{\Pi} \otimes \varepsilon)$ is proved in a similar way, and with this we have showed $(\overline{\Pi}, \Delta, \varepsilon)$ is a coalgebra in $C^{S_n}$.

It is enough to prove now that $\Delta$ and $\varepsilon$ are algebra morphisms. For $\Delta$ we have to show that the diagram

$$
\begin{array}{ccc}
\overline{\Pi} \otimes \overline{\Pi} & \xrightarrow{\Delta \otimes \Delta} & (\overline{\Pi} \otimes \overline{\Pi}) \otimes (\overline{\Pi} \otimes \overline{\Pi}) \\
\mu & \downarrow & \mu \\
\overline{\Pi} & \xrightarrow{\Delta} & \overline{\Pi} \otimes \overline{\Pi}
\end{array}
$$

commutes up to the relation $\hat{\mu}$, where $\hat{\mu}$ is the product in $\overline{\Pi} \otimes \overline{\Pi}$ and it is defined, as in (36), by the composite

$$
\hat{\mu} : (\overline{\Pi} \otimes \overline{\Pi}) \xrightarrow{A_{\overline{\Pi} \otimes \overline{\Pi}}} ((\overline{\Pi} \otimes \overline{\Pi}) \otimes (\overline{\Pi} \otimes \overline{\Pi})) \xrightarrow{A_{(\overline{\Pi} \otimes \overline{\Pi}) \otimes (\overline{\Pi} \otimes \overline{\Pi})}} \overline{\Pi} \otimes \overline{\Pi}
$$

The morphism $\text{Id}_\overline{\Pi} \otimes C_{\overline{\Pi} \otimes \overline{\Pi}} \otimes \text{Id}_\overline{\Pi}(v, w) : (\overline{\Pi} \otimes (\overline{\Pi} \otimes \overline{\Pi}))'(v) \rightarrow (\overline{\Pi} \otimes (\overline{\Pi} \otimes \overline{\Pi}))'(w)$ is related to $F_w(v, w) : (\overline{\Pi} \otimes (\overline{\Pi} \otimes \overline{\Pi}))(v) \rightarrow (\overline{\Pi} \otimes (\overline{\Pi} \otimes \overline{\Pi}))(w)$, which is represented by the following vertical arrow

$$
\begin{array}{c}
h(x_v)^* \otimes h(x_v) \otimes h(y_v)^* \otimes h(y_v) \otimes h(z_v)^* \otimes h(z_v) \otimes h(t_v)^* \otimes h(t_v) \\
F_w(v, w) = \delta_w^x \text{id}_{h(z_v)^* \otimes h(z_v)} \otimes (h(y_v)^*)^* \otimes h(y_v) \otimes \delta_w^x \text{id}_{h(t_v)^* \otimes h(t_v)} \\
h(x_w)^* \otimes h(x_w) \otimes h(z_w)^* \otimes h(y_w)^* \otimes h(y_w) \otimes h(t_w)^* \otimes h(t_w)
\end{array}
$$

since their codomains are related by associativity.

The morphism $(\mu \otimes \mu)(w, u) : (\overline{\Pi} \otimes \overline{\Pi}) \otimes (\overline{\Pi} \otimes \overline{\Pi}) \rightarrow \overline{\Pi} \otimes \overline{\Pi}$ is represented by the following vertical arrow

$$
\begin{array}{c}
h(x_w)^* \otimes h(x_w) \otimes h(z_w)^* \otimes h(y_w)^* \otimes h(y_w) \otimes h(t_w)^* \otimes h(t_w) \\
\delta_{x_w, x}(x_w, x_w) \Gamma_{h(z_w), h(x_w)} \otimes (\delta_{y_w, x}(x_w, y_w) \Gamma_{h(t_w), h(y_w)}) \\
h(x_u)^* \otimes h(x_u) \otimes h(y_u)^* \otimes h(y_u)
\end{array}
$$

It is not difficult to see that $\hat{\mu} \sim \sum_w (\mu \otimes \mu) \circ F_w$ and that this last morphism turns out to be equal to

$$
\begin{array}{c}
h(x_v)^* \otimes h(x_v) \otimes h(y_v)^* \otimes h(y_v) \otimes h(z_v)^* \otimes h(z_v) \otimes h(t_v)^* \otimes h(t_v) \\
\delta_{x_v, x}(x_v, x_v) \delta_{y_v, x}(x_v, y_v) \Gamma_{h(z_v), h(x_v)} \otimes (\delta_{y_v, x}(x_v, y_v) \Gamma_{h(t_v), h(y_v)}) (\text{id}_{h(x_v)^*} \otimes h(x_v)^* \otimes h(y_v)^* \otimes h(y_v)) \otimes (\text{id}_{h(t_v)^*} \otimes h(t_v)^* \otimes h(t_v)) \\
h(x_u)^* \otimes h(x_u) \otimes h(y_u)^* \otimes h(y_u)
\end{array}
$$
Hence $\tilde{\mu}(\Delta \otimes \Delta) = \sum_v G_v \circ (\Delta \otimes \Delta)$, where $G_v$ is the last vertical arrow. But $(\Delta \otimes \Delta)(p, v) : (\text{H} \otimes \text{H})(p) \rightarrow ((\text{H} \otimes \text{H}) \otimes (\text{H} \otimes \text{H}))(v)$ is given by

$$(\Delta \otimes \Delta)(p, v) = \delta_{x_v, x} \delta_{x_v, y_v} \delta_{y_v, z} \delta_{y_v, t_v} (\mu_{x_p} \otimes b_{x_p}) \otimes (\mu_{y_p} \otimes b_{y_p}) \otimes (\mu_{z_p} \otimes b_{z_p}) (\mu_{t_p} \otimes b_{t_p})$$

so the sum yields

$$M = \delta_{x_u, \chi(y, y)} \delta_{y_v, x_v} \delta_{y_v, y_v} \delta_{x_v, x_v} (\Gamma_{y_p, x_p} \otimes \Gamma_{y_v, x_v}) (\mu_{x_p} \otimes b_{x_p}) \otimes (\mu_{y_v} \otimes b_{y_v}) \otimes (\mu_{y_v} \otimes b_{y_v})$$

On the other hand, $(\mu \Delta)(p, u)$ is the sum over $v$ of the following composite

$$h(x_p) \otimes h(y_p) \otimes h(y_p) \xrightarrow{\delta_{x_u, x_v} \delta_{y_v, y_v} \delta_{y_v, y_v} \delta_{x_v, x_v} \Gamma_{y_p, x_p} \otimes \Gamma_{y_v, x_v}} h(x_v) \otimes h(x_v) \otimes h(x_v)$$

which is equal to

$$(\mu \Delta)(p, u) = \delta_{x_u, \chi(y, y)} \delta_{y_v, x_v} \delta_{y_v, y_v} \delta_{x_v, x_v} (\mu_{x_p} \otimes b_{x_p}) \otimes (\mu_{y_v} \otimes b_{y_v}) \otimes (\mu_{y_v} \otimes b_{y_v}) \Gamma_{h(y_v), h(x_v)}$$

In Figure 4 (at the end of the paper), taking $y = x_p$ and $x = y_p$, the picture on the upper left side represents $M$, while that on the lower right side represents $(\mu \Delta)(p, u)$. Hence both are equal and then $\mu \Delta = \tilde{\mu}(\Delta \otimes \Delta)$. Finally, we have to prove that $\varepsilon$ is an algebra morphism, that is, we have to prove that the diagram

$$\begin{array}{ccc}
\text{H} \otimes \text{H} & \xrightarrow{\varepsilon \otimes \varepsilon} & \text{I} \\
\mu \downarrow & & \\
\text{H} & \xrightarrow{\varepsilon} & \text{I}
\end{array}$$

commutes. We have

$$(\varepsilon \mu)(u, *) = \sum_{w} \varepsilon(\gamma(x_w, x_w), *) \circ \mu(u, w)$$

$$= \sum_{w} d_{h(x_w)} \circ (\delta_{x_w, \chi(y_u, x_u)} \Gamma_{h(y_u), h(x_u)})$$

$$= d_{h(\chi(y_u, x_u))} \Gamma_{h(y_u), h(x_u)}$$

On the other hand, $(\varepsilon \otimes \varepsilon)(u, *) = d_{h(x_u)} \otimes d_{h(y_u)}$.

Figure 5 (at the end of the paper), taking $x = y_u$ and $y = x_u$ as before, shows that these two morphisms are equal. Therefore $(\text{H}, \mu, \eta, \Delta, \varepsilon)$ is a bialgebra.

We shall now define the action of $\text{H}$ on the objects of $\mathcal{C}$. Take the point $x_0$ of $S_h$ such that $h(x_0) = \text{I}$ and define $j_V(x_0) = V$, for each $V$ object of $\mathcal{C}$. Define $T : \text{H} \otimes j_V \rightarrow j_V$, by

$$T(\gamma(x, x), x_0) = d_{h(x)} \otimes \text{id}_V : h^*(x) \otimes h(x) \otimes V \rightarrow V$$
where $x \in S_h$. It is not difficult to see that $T$ is indeed an action as we defined it before. The proof of that is similar (although easier and shorter) to the previous proofs and we omitted it.

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Figure 2: $\Gamma_{x,y,z}(\Gamma_{y,z} \otimes \text{id}_{x^*} \otimes x) = \Gamma_{x \otimes y,z}(\text{id}_{z^*} \otimes z \otimes \Gamma_{x,y})$
Figure 3: $(\text{id}_{h_1} \otimes b_{h_1} \otimes \text{id}_{h_1} \otimes \text{id}_{\overline{h_1} \otimes h_1})(\text{id}_{\overline{h_1} \otimes h_1} \otimes b_{h_1} \otimes \text{id}_{h_1}) = (\text{id}_{\overline{h_1} \otimes h_1} \otimes \text{id}_{\overline{h_1} \otimes h_1})(\text{id}_{\overline{h_1} \otimes h_1} \otimes b_{h_1} \otimes \text{id}_{h_1})$
Figure 4: \((\hat{\mu}(\Delta \otimes \Delta))(p, u) \doteq (\mu \Delta)(p, u)\)
Figure 5: $d_{h_i \otimes h_j} \Gamma_{h_i, h_j} = d_{h_j} \otimes d_{h_i}$
