Integrated Power and Subcarrier assignment by Cooperative Communications in LTE-A Networks

R.RajaKumar, R.Pandian, P.Indumathi, Jyotirmayee Subudhi

Abstract: The major enhancement of Long Term Evolution (LTE), namely, LTE-Advanced aims at larger capacity for the users. The specifications of LTE include, Downlink (DL) at 3Gbps, Uplink (UL) at 1.5 Gbps spectral efficiency at 30 bps/Hz and Type II relay station. The relay station and the eNodeB (eNB) make use of the same In-Band channel to maximize the available resources. This paper deals with optimal allocation of the Multicarrier scheme, namely, OFDM’s power and subcarrier to maximize DL communication efficiency of multiuser. Subcarriers are assigned to users having the highest channel quality and using water-filling strategy, the power is allocated. This allocation scheme is assessed against existing schemes and found to be effective. Thus our work aims at introducing power and resource allocation schemes in a LTE-A network.

Keywords—LTE-A, cooperative communication, spectrum efficiency, relay channel, water-filling algorithm, resource allocation, user fairness.

I. INTRODUCTION

Worldwide Third-Generation (3G) systems are employed to improve DL and UL transmissions. Anyway, because of the advancements in technology and higher demands of Service quality and Service experience requirement, future wireless systems face challenges such as higher data rates and efficient multimedia services. Hence 3GPP has released the LTE standard for wireless systems with the goal of increasing spectral efficiency. These choices make LTE-A support higher peak rates, throughput, coverage and reduced latencies improving QoE of user [3]. Driven by the user’s claiming extremely high rate connections, maximizing the network efficiency through resource and power allocation fuel slot of research work. By cooperative communications among base stations and the relays, cell-edge throughput and coverage are enhanced. A Type II relay is bidirectional and has transparency to every UE within its coverage, not needing any resource. This paper aims at partitioning. In our work we deal with Type II relays as they provide multipath diversity and UE gains. Among the cooperative protocols, Decode and Forward protocol has an advantage, as relay and source use the same channel increasing spectral efficiency.

- The optimal assignment of power between the eNB and the RS can improve data rate on every subcarrier. By this technique, the eNB and the RS cooperate so as to increase transmission efficiency.
- The optimal resource allocation scheme to increase the throughput on each subcarrier.

Our paper is arranged as below. The Section II provides system model. The maximization of the achievable rate of relay channel in a subcarrier is provided in Section III. The optimum power allocation among eNB and RS as well as the throughput maximization are contained in Section IV. Results of Simulation and concluding remarks are presented in Section V and Section VI respectively.

II. SYSTEM MODEL

The system model of our work is shown below. An eNB is deployed in every LTE-A cell at the center to provide service to many UEs. Relay Stations (RS) with minimal area of coverage are present near edges of the cell, so as to increase throughput of edge users. If there were no cooperation among neighbouring RS, users get service from their assigned eNB and RS if feasible. In DL, serving eNB sends signals to users. Suppose the user’s location is within range of coverage of RS, eNB reaches the user via both the direct link and the link via the relay, just as that is modeled by the cooperative relay channel with three nodes. DL communication is aided by the assigned RS as it captures signals transmitted from eNB and direct to user. Upon reception, RS retrieves the original message and sends again in its specific codes, while further transmission block occurs. From practical reality, Two-hop transmissions are dealt in this work. As hops increase, the complexity of decoding increases [19]. OFDM is the DL transmission scheme in LTE-A, where the available band B is orthogonally shared into n subchannels at various subcarriers, with each one having bandwidth 1/n of the total B Hz. The subcarriers can be assigned time to time so that the users can make use of both diversities of multiuser and frequency at a very fine level.

Revised Manuscript Received on August 30 2019.

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![System model](image)

**Fig. 1. System model**

Let the user set be denoted as $M = \{1, \ldots, m\}$ and subchannel sets be denoted as and $N = \{1, \ldots, n\}$ respectively. When a transmission block starts, eNB assigns the subchannels and the transmitted power is attained flexibly at evolved Node B and Relay Station maximizing efficiency of transmission. If all the wireless channels are assumed Gaussian, and the RS does not cover a user, the channel gain coefficient is $j_s h_i$. If Transmit and Receive end global Channel State Information (CSI), specific power and subchannel can be allocated by eNB optimally.

**Fig. 2. LTE-Advanced cellular network and Relay Stations.**

### III. SINGLE USER RELAY CHANNEL AND POWER ASSIGNMENT

The transmit-end node is assumed to have power restriction as

\[
\frac{1}{T} \sum_{i=1}^{T} x_s^2(t) \leq P_i,
\]

The relay node is considered to have power constraint as below while link $i$ is active.

\[
\frac{1}{T} \sum_{i=1}^{T} x_r^2(t) \leq P_i^{(i)}.
\]

Hence, the signals obtained by the nodes at relay and sink are, respectively

\[
y_s(t) = \frac{x_s(t)}{\sqrt{2} P_i} + z_s(t),
\]
\[
y_r(t) = \frac{x_r(t)}{\sqrt{2} P_i} + z_r(t) + z_d(t),
\]

where $z_s(t)$ and $z_d(t)$ are AWGN obtained at node of relay and at the node of sink $d$, with variance. Let $l_i$ represents OP distance. Path loss parameter $\eta$ has a range $2 - 8$, for all links. Let $P_i^{(i)}$ denote relay node’s power while transmitting in link $i$. While a block is transmitted, the node $r$ at relay just transmits the present message’s code with its maximum transmit power $P_i^{(i)}$. In the case of node at sources, the overall transmission power $P_i$ is divided into two parts, $P_i^{(i)}$ and $P_i^{(o)}$ for various purposes. $P_i^{(i)}$ is employed for the next message and $P_i^{(o)}$ is that of the relay for sending the present message to the sink. The two codes are combined and transmitted by $s$. Hence, the rate achieved for link $i$ with this relaying strategy is as below,

\[
R_i = \max_{P_i} \min \{I(X_s; Y_r | X_r), I(X_s, X_r; Y_d)\}
\]

**III. SINGLE USER RELAY CHANNEL AND POWER ASSIGNMENT**

\[
= \max_{P_i^{(i)}, P_i^{(o)}} \min \left\{ \frac{1}{2} \log \left( 1 + \frac{P_i^{(i)}}{P_i^{(o)} \sigma^2} \right), \frac{1}{2} \log \left( 1 + \frac{P_i^{(i)}}{P_i^{(o)} \sigma^2} + \frac{P_i^{(o)} \sigma^2}{P_i^{(o)} \sigma^2} \right) \right\}
\]

\[
= \max_{P_i^{(i)}, P_i^{(o)}} \min \left\{ \frac{1}{2} \log \left( 1 + \frac{P_i^{(i)}}{P_i^{(o)} \sigma^2} \right), \frac{1}{2} \log \left( 1 + \frac{P_i^{(i)}}{P_i^{(o)} \sigma^2} + \frac{P_i^{(o)} \sigma^2}{P_i^{(o)} \sigma^2} \right) \right\}
\]

We can observe that $R_i$ the rate achievable is dependent on $P_i^{(i)}$, $P_i^{(o)}$ and relay location. Our goal is to maximize throughput by joint assignment of power subject to constrained total power consumption and optimally locate the relay. This is the following optimization problem:

\[
\max \sum_{i=1}^{n} \frac{R_i}{P_i^{(i)} + P_i^{(o)} + P_i^{(f)} \leq P_i}, \forall i = 1, \ldots, n.
\]

here $P_i$ is the maximum power consumed while the active link is $i$. The first term $I(X_s; Y_r | X_r)$ is the decoding rate achieved by the relay, and the second term $I(X_s, X_r; Y_d)$ is the rate successfully decoded by the sink. Given a relay’s transmission power $P_i^{(i)}$, the maximum rate obtained is through the optimum assignment of $P_i^{(i)}$ and $P_i^{(o)}$ at the source. That is, the source node tries balancing the rates at which relay and sink decode and the ideal strategy is to have both the rates equal.

Based on the constraint rate, there are two choices available for the source code for power allocation:

- When the rate at which sink node decodes is the constraint, the node at source can enhance $P_i^{(o)}$ and decrease $P_i^{(i)}$ to make the two rates equal. This is synchronous case.
- When the rate at which relay decodes is the constraint, the source node sets $P_i^{(o)} = P$ and $P_i^{(i)} = 0$. While $P_i^{(o)} = 0$, the source and the relay communicate independently. This is asynchronous case.

In our optimization method, we will consider joint allocation of $P_i^{(i)}$, $P_i^{(o)}$, and $P_i^{(f)}$ for the two strategies.

#### A. Synchronous Scenario

In this case, as the sink receives the combined strength, initially we maximize the rate at the decoder of sink with fixed $P_i^{(i)}$, $P_i^{(o)} = P_i^{(2)} + P_i^{(f)}$. Later, we can distribute $P_i^{(i)}$ and $P_i^{(o)}$ with overall power constraint. The rate at which sink decodes is
I(X_a;X_r;Y_d) = \frac{1}{2} \log \left(1 + \frac{P_i(2)}{I_{slid} + I_{sir} + \sigma^2} \right).

from [19].

Thus, the optimal distribution between $P_i(2)$ and $P_r(i)$ is

$P_i(2) = \frac{I_{slid}}{I_{slid} + I_{sir} + \sigma^2}$

$P_r(i) = \frac{I_{sir}}{I_{slid} + I_{sir} + \sigma^2}$

The decoder rate at sink becomes:

$I(X_{s1};Y_r;Y_d) = \frac{1}{2} \log \left(1 + \frac{4P_0^{(i)} I_{l_0}^{(i)} I_{l_{slid}}}{I_{l_{sir}} I_{l_{sid}} + I_{l_{sir}} I_{l_{sid}} + \sigma^2} \right)$

Rewriting the optimization problem as

$\max \frac{1}{2} \log \left(1 + \frac{P_{l_0}^{(i)}}{\frac{I_{l_{sid}}}{I_{l_{sir}}}} \right)$

with constraint $P_{l_0}^{(i)} + P_{l_0}^{(i)} = P^{(i)}$

$P_{l_0}^{(i)} = \frac{4P_0^{(i)} I_{l_0}^{(i)} I_{l_{slid}}}{I_{l_{sir}} I_{l_{sid}} + I_{l_{sir}} I_{l_{sid}} + \sigma^2}$

If $l_{sid} \geq l_{sir}$ i.e., a relay is employed between the source and the sink to enable passing on the traffic, the optimal power allocation then is:

$P_i^{(1)} = \frac{4I_{l_0}^{(i)} I_{l_{slid}}}{I_{l_{sir}}} + P_{l_0}^{(i)}$

$P_r^{(i)} = \frac{4I_{l_0}^{(i)} I_{l_{slid}}}{I_{l_{sir}}} + P_{l_0}^{(i)}$

from [19].

The maximum rate achieved is:

$R^{\text{max}} = \frac{1}{2} \log \left(1 + \frac{4I_{l_0}^{(i)} I_{l_{slid}}}{I_{l_{sir}} + \sigma^2} \right)$

If $\beta(k,l) \in [0, 1]$ is the parameter that the transmit end employs, to finetune the part of power transmitted that is used to cooperate with the relay so as to attain the highest rate and $\beta(k,l) = 1 - \beta(k,l)$. We can also write

$R^{(k,l)} = \frac{1}{n} \log \left(1 + \frac{|h_{sr}^{(k,l)}|^2 |h_{sr}^{(k,l)}|^2 + |h_{sr}^{(k,l)}|^2 |h_{sr}^{(k,l)}|^2}{|h_{sr}^{(k,l)}|^2 + |h_{sr}^{(k,l)}|^2} \right)$

from [19].

B. Asynchronous Scenario

In asynchronous ($l_{sid} \geq l_{sir}$) case, $P_i^{(2)} = 0$ and $P_r^{(2)} = P_k$ the optimization problem becomes

$\max \frac{1}{2} \log \left(1 + \frac{P_i(1)}{I_{slid} + I_{sir} + \sigma^2} \right)$

subject to $P_i + P_r^{(i)} = P^{(i)}$.

Similarly, the obtained rate is maximum when,

$P_i^{(2)} = \frac{I_{slid}}{I_{slid} + I_{sir} + \sigma^2}$

$P_r^{(i)} = \frac{I_{sir}}{I_{slid} + I_{sir} + \sigma^2}$

From [20]

Hence, the optimal power sharing here is

$P_i = \frac{I_{slid}}{I_{slid} + I_{sir} + \sigma^2} - \frac{I_{slid} (I_{slid} + I_{sir} + \sigma^2)}{I_{slid} + I_{sir} + \sigma^2}$

$P_r^{(i)} = \frac{I_{sir} I_{sir} + I_{sir} + \sigma^2}{I_{slid} + I_{sir} + \sigma^2} - \frac{I_{slid} (I_{slid} + I_{sir} + \sigma^2)}{I_{slid} + I_{sir} + \sigma^2}$

The highest achievable rate is:

$R_s^{(\text{async})} = \frac{1}{2} \log \left(1 + \frac{P_i}{I_{slid} + I_{sir} + \sigma^2} \right)$

Or

$R^{(k,l)} = \frac{1}{n} \log \left(1 + \frac{|h_{sr}^{(k,l)}|^2 + |h_{sr}^{(k,l)}|^2}{|h_{sr}^{(k,l)}|^2 + |h_{sr}^{(k,l)}|^2} \right)$

from [20]

C. Non-cooperative case

If $l_{sid} < l_{sir}$ then $P_i^{(2)} = 0$ and implies that the relay decoder is the bottleneck in all scenarios of power allocation. Thus, the distribution in optimum manner can be obtained when $P_i^{(1)} = P_0$. This is the case wherein the relay is not utilized for source-sink transmission and the maximum rate is given by

$R^{\text{sync}} = \frac{1}{2} \log \left(1 + \frac{P_i}{I_{slid} + \sigma^2} \right)$

Or

$R^{(k,l)} = \frac{1}{n} \log \left(1 + \frac{|h_{sr}^{(k,l)}|^2 + |h_{sr}^{(k,l)}|^2}{|h_{sr}^{(k,l)}|^2 + |h_{sr}^{(k,l)}|^2} \right)$

D. Equal power allocation

In this scheme the eNB sends with the same power as the RS.

IV. JOINT POWER AND SUBCHANNEL ALLOCATION FOR THROUGHPUT MAXIMIZATION

A. Existing Scheme – Scheme 1

Statement: To increase the data rate of a given subcarrier in a DL multiuser OFDM scheme, is that the subcarrier should be allotted to a single user with the highest channel gain in that subcarrier. Hence, a subcarrier is used by only one user and never by multiple users in a given time. After the allocation of subcarriers, the total transmit power is equally assigned to the subcarriers.

$S_m = S/M$ for $m = 1, 2, \ldots, M$. Here $S$ is the power transmitted and $M$ is the total number of sub channels.

B. Existing Scheme – Scheme 2

The power and subchannels are equally distributed for all users, unmindful of each user’s channel conditions. There are $K$ users overall consuming $N$ sub channels,
while transmitted power is constrained to P.

C. Algorithm for Optimal Resource Assignment

The goal here is to increase the overall throughput.

\[ \max_{p^{(k)}(l), p^{(l)}} \sum_{k \in M, l \in N} R^{(k,l)} \]

given, \( p^{(k)}(l) \leq P_{\text{total}} \), \( p^{(k)}(l) \geq 0 \), for all \( k, l \), \( \rho^{(l)} \in \{0,1\} \), for all \( k, l \)

Optimally allocating the resources jointly, amount to increase the overall throughput by assigning the subcarrier to the user with best channel and allotting power by water-filling strategy.

- Given, \( l = 1, \ldots, n \), find a \( k(l) \) fulfilling the condition \( H^{(k,l)} \geq H^{(k,l)} \) for every \( k \in M \). Allot subchannel \( l \) to specific user \( k(l) \), i.e., set \( p^{(k(l),l)} = 1 \).
- Assign \( p^{(k(l),l)} = (\lambda - 1 / H^{(k,l),l})^+ \) as the power transmitted for the user \( k(l) \) in subchannel \( l \). \( \lambda \) is the level of water-filling process that is chosen to fulfill the total power restriction.

D. Water filling process.

To sum it up in a nutshell, as power of signal is raised from zero, we allocate the power for the channel with lowermost noise.

V. NUMERICAL RESULTS

To understand our joint resource assignment scheme, the results of simulation are stated as two parts.

A. Power Assignment and Single-User Relay Channel

The rates attainable for an individual user are obtained by optimally dividing power for both synchronous and asynchronous cases. We compare it with equal power division where both evolved Node B and the Relay Station have equal power. To understand the improvement due to employing relay station, we keep capacity of transmission without cooperation as a baseline to compare, as shown in Fig. 3. We deduct that cooperative transmission with relay outperforms the relay-less noncooperative strategy through achievable rate of data. Also, higher rate can be achieved in synchronous scenario than asynchronous scenario most often, since the former one makes use of the cooperation between the eNB and RS while coding. It is also observed that equal power division is inferior to our scheme of optimum power division.

B. Resource Assignment for Throughput Maximization

Fig. 4 compares the throughput of the three schemes of synchronous scenario and Fig. 5 shows the behaviour when asynchronous relay is considered. From these figures, it can be observed that the highest throughput is attained by optimal resource assignment in both synchronous and asynchronous cases. Moreover, scheme 1 beats scheme 2, as it uses the channel variance. Also, we can infer from both the figures, that the achievable throughput of the optimal resource allocation and the one attained by scheme 1 increases as users increase.

VI. CONCLUSION

We have studied the DL resource allocation for maximizing the communication efficiency in an LTE A system by deploying relays. Contrary to just amplify and forward the signal received, cooperation is required by relays to achieve maximum data rates. Same channel occupancy of source and relays is the property of In-Band Type II decode-and-forward relays employed in our work. An optimum joint assignment of subcarriers and power is employed here for increasing achieved throughput. Simulation results confirm that our algorithm increases the minimum rate of each user. Relay installation with cooperation for DL-COMP can be a future work.
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