Forest-Based Translation

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Abstract

Among syntax-based translation models, the tree-based approach, which takes as input a parse tree of the source sentence, is a promising direction being faster and simpler than its string-based counterpart. However, current tree-based systems suffer from a major drawback: they only use the 1-best parse to direct the translation, which potentially introduces translation mistakes due to parsing errors. We propose a forest-based approach that translates a packed forest of exponentially many parses, which encodes many more alternatives than standard n-best lists. Large-scale experiments show an absolute improvement of 1.7 BLEU points over the 1-best baseline. This result is also 0.8 points higher than decoding with 30-best parses, and takes even less time.

1 Introduction

Syntax-based machine translation has witnessed promising improvements in recent years. Depending on the type of input, these efforts can be divided into two broad categories: the string-based systems whose input is a string to be simultaneously parsed and translated by a synchronous grammar (Wu, 1997; Chiang, 2005; Galley et al., 2006), and the tree-based systems whose input is already a parse tree to be directly converted into a target tree or string (Lin, 2004; Ding and Palmer, 2005; Quirk et al., 2005; Liu et al., 2006; Huang et al., 2006). Compared with their string-based counterparts, tree-based systems offer some attractive features: they are much faster in decoding (linear time vs. cubic time, see (Huang et al., 2006)), do not require a binary-branching grammar as in string-based models (Zhang et al., 2006), and can have separate grammars for parsing and translation, say, a context-free grammar for the former and a tree substitution grammar for the latter (Huang et al., 2006). However, despite these advantages, current tree-based systems suffer from a major drawback: they only use the 1-best parse tree to direct the translation, which potentially introduces translation mistakes due to parsing errors (Quirk and Corston-Oliver, 2006). This situation becomes worse with resource-poor source languages without enough Treebank data to train a high-accuracy parser.

One obvious solution to this problem is to take as input k-best parses, instead of a single tree. This k-best list postpones some disambiguation to the decoder, which may recover from parsing errors by getting a better translation from a non 1-best parse. However, a k-best list, with its limited scope, often has too few variations and too many redundancies; for example, a 50-best list typically encodes a combination of 5 or 6 binary ambiguities (since $2^5 < 50 < 2^6$), and many subtrees are repeated across different parses (Huang, 2008). It is thus inefficient either to decode separately with each of these very similar trees. Longer sentences will also aggravate this situation as the number of parses grows exponentially with the sentence length.

We instead propose a new approach, forest-based translation (Section 3), where the decoder translates a packed forest of exponentially many parses,\footnote{There has been some confusion in the MT literature regarding the term forest: the word “forest” in “forest-to-string rules” has several different meanings depending on the context.}
which compactly encodes many more alternatives than $k$-best parses. This scheme can be seen as a compromise between the string-based and tree-based methods, while combining the advantages of both: decoding is still fast, yet does not commit to a single parse. Large-scale experiments (Section 4) show an improvement of 1.7 BLEU points over the 1-best baseline, which is also 0.8 points higher than decoding with 30-best trees, and takes even less time thanks to the sharing of common subtrees.

2 Tree-based systems

Current tree-based systems perform translation in two separate steps: parsing and decoding. A parser first parses the source language input into a 1-best tree $T$, and the decoder then searches for the best derivation (a sequence of translation steps) $d^*$ that converts source tree $T$ into a target-language string among all possible derivations $D$:

$$d^* = \arg \max_{d \in D} P(d | T). \quad (1)$$

We will now proceed with a running example translating from Chinese to English:

(2) 布什与沙龙举行了会谈

Bush with/and Sharon held pass. talk

“Bush held a talk with Sharon”

Figure 2 shows how this process works. The Chinese sentence (a) is first parsed into tree (b), which will be converted into an English string in 5 steps. First, at the root node, we apply rule $r_1$ preserving top-level word-order between English and Chinese,

$$(r_1) \ \text{IP}(x_1:NPB \ x_2:VP) \rightarrow x_1 \ x_2$$

(Liu et al., 2007) was a misnomer which actually refers to a set of several unrelated subtrees over disjoint spans, and should not be confused with the standard concept of packed forest.

Figure 2: An example derivation of tree-to-string translation. Shaded regions denote parts of the tree that is pattern-matched with the rule being applied.

which results in two unfinished subtrees in (c). Then rule $r_2$ grabs the $B\text{ Bush}$ subtree and transliterate it

$$(r_2) \ \text{NPB}(\text{NR}(B\text{ Bush})) \rightarrow \text{Bush}.$$
More formally, a (tree-to-string) translation rule (Huang et al., 2006) is a tuple \( \langle t, s, \phi \rangle \), where \( t \) is the source-side tree, whose internal nodes are labeled by nonterminal symbols in \( \mathcal{N} \), and whose frontier nodes are labeled by source-side terminals in \( \Sigma \) or variables from a set \( \mathcal{X} = \{ x_1, x_2, \ldots \} \); \( s \in (\mathcal{X} \cup \Delta)^* \) is the target-side string where \( \Delta \) is the target language terminal set; and \( \phi \) is a mapping from \( \mathcal{X} \) to nonterminals in \( \mathcal{N} \). Each variable \( x_i \in \mathcal{X} \) occurs exactly once in \( t \) and exactly once in \( s \). We denote \( \mathcal{R} \) to be the translation rule set. A similar formalism appears once in (Galley et al., 2004). 

In the reverse direction of the original string-to-tree transducer rules defined by Galley et al. (2004), the depth-first traversal procedure in tree-based decoding (Figure 2), we visit in top-down order each node \( v \) in the transducer rules defined by Galley et al. (2004). These rules are in the reverse direction of the original string-to-tree transducer rules defined by Galley et al. (2004).

Finally, from step (d) we apply rules \( r_4 \) and \( r_5 \)

\[\begin{align*}
(r_4) \quad & \text{NPB(NN(hu`ıt´an))} \rightarrow \text{a talk} \\
(r_5) \quad & \text{NPB(NR(Sh¯al´ong))} \rightarrow \text{Sharon}
\end{align*}\]

which perform phrasal translations for the two remaining subtrees, respectively, and get the Chinese translation in (e).

3 Forest-based translation

We now extend the tree-based idea from the previous section to the case of forest-based translation. Again, there are two steps, parsing and decoding. In the former, a (modified) parser will parse the input sentence and output a packed forest (Section 3.1) rather than just the 1-best tree. Such a forest is usually huge in size, so we use the forest pruning algorithm (Section 3.4) to reduce it to a reasonable size. The pruned parse forest will then be used to direct the translation.

In the decoding step, we first convert the parse forest into a translation forest using the translation rule set, by similar techniques of pattern-matching from tree-based decoding (Section 3.2). Then the decoder searches for the best derivation on the translation forest and outputs the target string (Section 3.3).

3.1 Parse Forest

Informally, a packed parse forest, or forest in short, is a compact representation of all the derivations (i.e., parse trees) for a given sentence under a context-free grammar (Billot and Lang, 1989). For example, consider the Chinese sentence in Example (2) above, which has (at least) two readings depending on the part-of-speech of the word yū, which can be either a preposition (P “with”) or a conjunction (CC “and”). The parse tree for the preposition case is shown in Figure 2(b) as the 1-best parse, while for the conjunction case, the two proper nouns (Bush and Shālōng) are combined to form a coordinated NP

\[
\begin{array}{cccc}
\text{NPB}_{0,1} & \text{CC}_{1,2} & \text{NPB}_{2,3} \\
\text{NP}_{0,3} \\
\end{array}
\]

(*)

which functions as the subject of the sentence. In this case the Chinese sentence is translated into

(3) “[Bush and Sharon] held a talk”.

Shown in Figure 3(a), these two parse trees can be represented as a single forest by sharing common subtrees such as \( \text{NPB}_{0,1} \) and \( \text{VPB}_{3,6} \). Such a forest has a structure of a hypergraph (Klein and Manning, 2001; Huang and Chiang, 2005), where items like \( \text{NP}_{0,3} \) are called nodes, and deductive steps like (*) correspond to hyperedges.

More formally, a forest is a pair \( \langle V, E \rangle \), where \( V \) is the set of nodes, and \( E \) the set of hyperedges. For a given sentence \( w_{1:l} = w_1 \ldots w_l \), each node \( v \in V \) is in the form of \( X_{i:j} \), which denotes the recognition of nonterminal \( X \) spanning the substring from positions \( i \) through \( j \) (that is, \( w_{i+1} \ldots w_j \)). Each hyperedge \( e \in E \) is a pair \( (\text{tails}(e), \text{head}(e)) \), where \( \text{head}(e) \in V \) is the consequent node in the deductive step, and \( \text{tails}(e) \in V^* \) is the list of antecedent nodes. For example, the hyperedge for deduction (*) is notated:

\[
\langle \langle \text{NPB}_{0,1}, \text{CC}_{1,2}, \text{NPB}_{2,3} \rangle, \text{NP}_{0,3} \rangle.
\]

There is also a distinguished root node TOP in each forest, denoting the goal item in parsing, which is simply \( S_{0,l} \) where \( S \) is the start symbol and \( l \) is the sentence length.

3.2 Translation Forest

Given a parse forest and a translation rule set \( \mathcal{R} \), we can generate a translation forest which has a similar hypergraph structure. Basically, just as the depth-first traversal procedure in tree-based decoding (Figure 2), we visit in top-down order each node \( v \) in the
parse forest, and try to pattern-match each translation rule $r$ against the local sub-forest under node $v$. For example, in Figure 3(a), at node $VP_{1,6}$, two rules $r_3$ and $r_7$ both matches the local subforest, and will thus generate two translation hyperedges $e_3$ and $e_4$ (see Figure 3(b-c)).

More formally, we define a function $\text{match}(r, v)$ which attempts to pattern-match rule $r$ at node $v$ in the parse forest, and in case of success, returns a list of descendent nodes of $v$ that are matched to the variables in $r$, or returns an empty list if the match fails. Note that this procedure is recursive and may

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Figure 3: (a) the parse forest of the example sentence; solid hyperedges denote the 1-best parse in Figure 2(b) while dashed hyperedges denote the alternative parse due to Deduction (*). (b) the corresponding translation forest after applying the translation rules (lexical rules not shown); the derivation shown in bold solid lines ($e_1$ and $e_3$) corresponds to the derivation in Figure 2; the one shown in dashed lines ($e_2$, $e_5$, and $e_6$) uses the alternative parse and corresponds to the translation in Example (3). (c) the correspondence between translation hyperedges and translation rules.

| translation hyperedge | translation rule |
|-----------------------|------------------|
| $e_1$                 | $r_1$            |
| $e_2$                 | $r_6$            |
| $e_3$                 | $r_3$            |
| $e_4$                 | $r_7$            |
| $e_5$                 | $r_8$            |
| $e_6$                 | $r_9$            |

$\text{IP}(x_1:\text{NP} x_2:\text{VP}) \rightarrow x_1 x_2$

$\text{IP}(x_1:\text{NP} x_2:\text{VPB}) \rightarrow x_1 x_2$

$\text{VP}(\text{PP}(\text{P}(\text{yú}) x_1:\text{NP} x_2:\text{VPB}) \text{VPB}(\text{V}(\text{jùxíng}) \text{AS}(\text{le}) x_2:\text{NPB})) \rightarrow \text{held } x_2 \text{ with } x_1$

$\text{VP}(\text{PP}(\text{P}(\text{yú}) x_1:\text{NP} x_2:\text{VPB}) \text{VPB}(\text{V}(\text{jùxíng}) \text{AS}(\text{le}) x_2:\text{NPB}) \rightarrow \text{held } x_1$

$\text{NP}(x_1:\text{NP} x_2:\text{NPB} \text{CC}(\text{yú}) x_2:\text{NPB}) \rightarrow x_1 \text{ and } x_2$

$\text{VPB}(\text{V}(\text{jùxíng}) \text{AS}(\text{le}) x_1:\text{NPB}) \rightarrow \text{held } x_1$
involve multiple parse hyperedges. For example,

\[
mATCH(r_3, VP_{1,6}) = (NPB_{2,3}, NPB_{5,6}),
\]

which covers three parse hyperedges, while nodes in gray do not pattern-match any rule (although they are involved in the matching of other nodes, where they match interior nodes of the source-side tree fragments in a rule). We can thus construct a translation hyperedge from \( \text{match}(r, v) \) to \( v \) for each node \( v \) and rule \( r \). In addition, we also need to keep track of the target string \( s(r) \) specified by rule \( r \), which includes target-language terminals and variables. For example, \( s(r_3) = \text{“held } x_2 \text{ with } x_1” \). The subtranslations of the matched variable nodes will be substituted for the variables in \( s(r) \) to get a complete translation for node \( v \). So a translation hyperedge \( e \) is a triple \( \langle \text{tails}(e), \text{head}(e), s \rangle \) where \( s \) is the target string from the rule, for example,

\[
e_3 = \langle (NPB_{2,3}, NPB_{5,6}), VP_{1,6}, \text{“held } x_2 \text{ with } x_1” \rangle.
\]

This procedure is summarized in Pseudocode 1.

**Pseudocode 1** The conversion algorithm.

1. **Input**: parse forest \( H_p \) and rule set \( R \)
2. **Output**: translation forest \( H_t \)
3. **for** each node \( v \in V_p \) in top-down order **do**
4. **for** each translation rule \( r \in R \) **do**
5. \( \text{vars} \leftarrow \text{match}(r, v) \) \( \triangleright \) variables
6. **if** \( \text{vars} \) is not empty **then**
7. \( e \leftarrow \langle \text{vars}, v, s(r) \rangle \)
8. add translation hyperedge \( e \) to \( H_t \)

For \( k \)-best search after getting 1-best derivation, we use the lazy Algorithm 3 of Huang and Chiang (2005) that works backwards from the root node, incrementally computing the second, third, through the \( k \)th best alternatives. However, this time we work on a finer-grained forest, called translation+LM forest, resulting from the intersection of the translation forest and the LM, with its nodes being the +LM items during cube pruning. Although this new forest is prohibitively large, Algorithm 3 is very efficient with minimal overhead on top of 1-best.

### 3.4 Forest Pruning Algorithm

We use the pruning algorithm of (Jonathan Graehl, p.c.; Huang, 2008) that is very similar to the method based on marginal probability (Charniak and Johnson, 2005), except that it prunes hyperedges as well as nodes. Basically, we use an Inside-Outside algorithm to compute the Viterbi inside cost \( \beta(v) \) and the Viterbi outside cost \( \alpha(v) \) for each node \( v \), and then compute the merit \( \alpha \beta(e) \) for each hyperedge:

\[
\alpha \beta(e) = \alpha(\text{head}(e)) + \sum_{u_i \in \text{tails}(e)} \beta(u_i) \quad (4)
\]

Intuitively, this merit is the cost of the best derivation that traverses \( e \), and the difference \( \delta(e) = \alpha \beta(e) - \beta(\text{TOP}) \) can be seen as the distance away from the globally best derivation. We prune away a hyperedge \( e \) if \( \delta(e) > p \) for a threshold \( p \). Nodes with all incoming hyperedges pruned are also pruned.

### 4 Experiments

We can extend the simple model in Equation 1 to a log-linear one (Liu et al., 2006; Huang et al., 2006):

\[
d^* = \arg \max_{d \in D} P(d \mid T)^{\lambda_0} \cdot e^{\lambda_1 |d|} \cdot P_{\text{lm}}(s)^{\lambda_2} \cdot e^{\lambda_3 |s|}
\]

where \( T \) is the 1-best parse, \( e^{\lambda_1 |d|} \) is the penalty term on the number of rules in a derivation, \( P_{\text{lm}}(s)^{\lambda_2} \) is the language model and \( e^{\lambda_3 |s|} \) is the length penalty term up the computation. An +LM item of node \( v \) has the form \( (v^{a \cdot b}, \alpha, \beta) \), where \( a \) and \( b \) are the target-language boundary words. For example, \((VP_{1,6}^{\text{held} \cdot \text{Sharon}})\) is an +LM item with its translation starting with “held” and ending with “Sharon”. This scheme can be easily extended to work with a general \( n \)-gram by storing \( n - 1 \) words at both ends (Chiang, 2007).
on target translation. The derivation probability conditioned on 1-best tree, \( P(d \mid T) \), should now be replaced by \( P(d \mid H_p) \) where \( H_p \) is the parse forest, which decomposes into the product of probabilities of translation rules \( r \in d \):

\[
P(d \mid H_p) = \prod_{r \in d} P(r) \tag{6}
\]

where each \( P(r) \) is the product of five probabilities:

\[
P(r) = P(t \mid s)^{\lambda_4} \cdot P_{\text{lex}}(t \mid s)^{\lambda_5} \cdot P(s \mid t)^{\lambda_6} \cdot P_{\text{lex}}(s \mid t)^{\lambda_7} \cdot P(t \mid H_p)^{\lambda_8}. \tag{7}
\]

Here \( t \) and \( s \) are the source-side tree and target-side string of rule \( r \), respectively, \( P(t \mid s) \) and \( P(s \mid t) \) are the two translation probabilities, and \( P_{\text{lex}}(\cdot) \) are the lexical probabilities. The only extra term in forest-based decoding is \( P(t \mid H_p) \) denoting the source side parsing probability of the current translation rule \( r \) in the parse forest, which is the product of probabilities of each parse hyperedge \( e_p \) covered in the pattern-match of \( t \) against \( H_p \) (which can be recorded at conversion time):

\[
P(t \mid H_p) = \prod_{e_p \in H_p, e_p \text{ covered by } t} P(e_p). \tag{8}
\]

### 4.1 Data preparation

Our experiments are on Chinese-to-English translation, and we use the Chinese parser of Xiong et al. (2005) to parse the source side of the bitext. Following Huang (2008), we modify the parser to output a packed forest for each sentence.

Our training corpus consists of 31,011 sentence pairs with 0.8M Chinese words and 0.9M English words. We first word-align them by GIZA++ refined by “diagand” from Koehn et al. (2003), and apply the tree-to-string rule extraction algorithm (Galley et al., 2006; Liu et al., 2006), which resulted in 346K translation rules. Note that our rule extraction is still done on 1-best parses, while decoding is on \( k \)-best parses or packed forests. We also use the SRI Language Modeling Toolkit (Stolcke, 2002) to train a trigram language model with Kneser-Ney smoothing on the English side of the bitext.

We use the 2002 NIST MT Evaluation test set as our development set (878 sentences) and the 2005 NIST MT Evaluation test set as our test set (1082 sentences), with on average 28.28 and 26.31 words per sentence, respectively. We evaluate the translation quality using the case-sensitive BLEU-4 metric (Papineni et al., 2002). We use the standard minimum error-rate training (Och, 2003) to tune the feature weights to maximize the system’s BLEU score on the dev set. On dev and test sets, we prune the Chinese parse forests by the forest pruning algorithm in Section 3.4 with a threshold of \( p = 12 \), and then convert them into translation forests using the algorithm in Section 3.2. To increase the coverage of the rule set, we also introduce a default translation hyperedge for each parse hyperedge by monotonically translating each tail node, so that we can always at least get a complete translation in the end.

### 4.2 Results

The BLEU score of the baseline 1-best decoding is 0.2325, which is consistent with the result of 0.2302 in (Liu et al., 2007) on the same training, development and test sets, and with the same rule extraction procedure. The corresponding BLEU score of Pharaoh (Koehn, 2004) is 0.2182 on this dataset.

Figure 4 compares forest decoding with decoding on \( k \)-best trees in terms of speed and quality. Using more than one parse tree apparently improves the BLEU score, but at the cost of much slower decoding, since each of the top-\( k \) trees has to be decoded individually although they share many common sub-trees. Forest decoding, by contrast, is much faster.
and produces consistently better BLEU scores. With pruning threshold $p = 12$, it achieved a BLEU score of $0.2485$, which is an absolute improvement of $1.6\%$ points over the 1-best baseline, and is statistically significant using the sign-test of Collins et al. (2005) ($p < 0.01$).

We also investigate the question of how often the $i$th-best parse tree is picked to direct the translation ($i = 1, 2, \ldots$), in both $k$-best and forest decoding schemes. A packed forest can be roughly viewed as a (virtual) $\infty$-best list, and we can thus ask how often is a parse beyond top-$k$ used by a forest, which relates to the fundamental limitation of $k$-best lists. Figure 5 shows that, the 1-best parse is still preferred $25\%$ of the time among 30-best trees, and $23\%$ of the time by the forest decoder. These ratios decrease dramatically as $i$ increases, but the forest curve has a much longer tail in large $i$. Indeed, $40\%$ of the trees preferred by a forest is beyond top-30, $32\%$ is beyond top-100, and even $20\%$ beyond top-1000. This confirms the fact that we need exponentially large $k$-best lists with the explosion of alternatives, whereas a forest can encode these information compactly.

### 4.3 Scaling to large data

We also conduct experiments on a larger dataset, which contains 2.2M training sentence pairs. Besides the trigram language model trained on the English side of these bitext, we also use another trigram model trained on the first 1/3 of the Xinhua portion of Gigaword corpus. The two LMs have distinct weights tuned by minimum error rate training. The dev and test sets remain the same as above.

Furthermore, we also make use of bilingual phrases to improve the coverage of the ruleset. Following Liu et al. (2006), we prepare a phrase-table from a phrase-extractor, e.g. Pharaoh, and at decoding time, for each node, we construct on-the-fly flat translation rules from phrases that match the source-side span of the node. These phrases are called syntactic phrases which are consistent with syntactic constituents (Chiang, 2005), and have been shown to be helpful in tree-based systems (Galley et al., 2006; Liu et al., 2006).

The final results are shown in Table 1, where TR denotes translation rule only, and TR+BP denotes the inclusion of bilingual phrases. The BLEU score of forest decoder with TR is $0.2839$, which is a $1.7\%$ points improvement over the 1-best baseline, and this difference is statistically significant ($p < 0.01$). Using bilingual phrases further improves the BLEU score by $3.1\%$ points, which is $2.1\%$ points higher than the respective 1-best baseline. We suspect this larger improvement is due to the alternative constituents in the forest, which activates many syntactic phrases suppressed by the 1-best parse.

### 5 Conclusion and future work

We have presented a novel forest-based translation approach which uses a packed forest rather than the 1-best parse tree (or $k$-best parse trees) to direct the translation. Forest provides a compact data-structure for efficient handling of exponentially many tree structures, and is shown to be a promising direction with state-of-the-art translation results and reasonable decoding speed. This work can thus be viewed as a compromise between string-based and tree-based paradigms, with a good trade-off between speed and accuracy. For future work, we would like to use packed forests not only in decoding, but also for translation rule extraction during training.

| approach \ ruleset | TR     | TR+BP  |
|-------------------|--------|--------|
| 1-best tree       | 0.2666 | 0.2939 |
| 30-best trees     | 0.2755 | 0.3084 |
| forest ($p = 12$) | **0.2839** | **0.3149** |

Table 1: BLEU score results from training on large data.
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