Multiple error correction towards optimisation of energy in sensor network

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ABSTRACT

Hybrid ARQ (HARQ) is among the optimum error controls implemented in Wireless Sensor Network as it reduces the overhead from retransmission and error correcting codes. The advancement in WSN includes the usage of high number of nodes and the increase in traffic with large data transmitted among the nodes had concerned the need for a new approach in error control algorithm. This paper proposed the multiple error correction based on HARQ process to aid the changes in channel with proper error correction assignment towards optimising the performances of WSN in terms of bit error rates, remaining energy, and latency for different types of congestion and channel conditions. In this study, we have developed the channel adaptation algorithm that can adapt to sudden changes and demonstrated the optimal error correcting codes as well as adjustment on the transmit power for the given channel condition and congestion presented. From the result analysed, the optimisation between the remaining energy and Bit Error rates happened on the basis of adapting to these different channel condition and congestion to minimize redundancies appended. From the result obtained, we concluded that by using multiple error correction algorithm with the aid of adjustment on the transmit power, the remaining energy can be optimised together with Bit Error rates and the excessive redundancies can be reduced.

Keywords:
HARQ
Kalman filter
Remaining energy
SNR channel adaptation
Wireless sensor network

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1. INTRODUCTION

The Wireless Sensor Network (WSN) is very crucial in monitoring field such as habitat monitoring, environmental, agricultural, military, and tracking field. Recently, there are some emerging applications of WSN in big data [1],[2] and internet of things (IoT) [3],[4]. The existing technology of WSN critically benefits in terms of cost, scalability, and also provide supports towards human-work constrained when monitoring dangerous places such as natural disasters and unfriendly environments [5].

As WSN is energy-constrained and error-prone, researchers have established many methods and ways to overcome these problems. Researches back then provided the method to reduce energy consumption through calibrating or adjusting the transmission power such as by estimating the Signal to Noise Ratio (SNR) [6] and Received Signal Strength Indicator RSSI [7] using Kalman Filter (KF) in order to adjust the transmission power. The important aspect of high error rates in a network cannot be cast aside although minimizing the energy usage using transmission power control (TPC) is an effective method to maintain the lifetime as high error rates cause retransmission to flood the network and this consumes more energy. Thus, the problem of high energy usage needs to be tackled alongside with high error rates. In addition, some
techniques in implementing error correction schemes such as HARQ indirectly reduce error rates and possibly reduce the energy consumption by reducing the flooding of retransmission in a congested network. There are related studies on the effect of HARQ with BCH in Coded Division Multiple Access (CDMA) WSN towards the error rates and energy consumption [8], [9].

The authors reported on high Bit Error Rates (BER) at low hop count. Thus, the architecture or node deployment contributes to the increase in error rates or packet corruption in the presence of interferences and noises particularly when the nodes are located within the transmission range of other nodes. Practically, the nodes are usually deployed in a non-uniform manner following the geographical structure. Thus, there can be variations of node densities in one monitoring area. Moreover, apart from interferences and signal fading, the channel condition might change over time due to reflections and refractions of signals and also might be contributed from the hardware itself. Thus, the implementation of existing error correction schemes might not be reliable as different levels of congestions might present in one monitoring area with sudden changes of interferences and noise over time. The work of [10] introduced the multi-coding schemes for WSN where different coding schemes were implemented at the nodes and the sink. Based from this literature, it was demonstrated that the multi-coding schemes increased the lifetime of a network. However, further studies on the optimal codeword length and error correcting capability that corresponds to the presence of noises are needed in order to fully optimize the energy usage and error rates in a network.

Different coding schemes and error correcting capability will append different number of redundancies to the transmitted bits. Based on our previous studies [11], [12], [13], we demonstrated that high redundancy consumed too much energy in decoding the transmitted data. Meanwhile, a network with bad condition and high error rates might not able to handle large number of erroneous bits due to the relatively low error correcting capability. Thus, in this paper, we proposed the approach to optimize the energy usage with BER by modifying the Hybrid ARQ (HARQ) process aided with multi-coding schemes and power control that can adapt to the changes in channel condition that were estimated as SNR using KF. Our proposed algorithm merits on its capability to adapt to the changes in SNR in which every SNR range describe different error correction used and transmit power calibration to minimise the higher redundancies when these excessive redundancies are not needed in better SNR condition and vice versa. Transmit power calibration aided to reduce the interferences and noises in noisy channel when SNR is low as the ratio of noises will rise the BER sharply. We presented the comparisons between our proposed algorithm towards the increase in performances with the existing error correction used in Section 3 in which we able to reduce high BER in low SNR condition and maintain remaining energy in high SNR condition.

2. RESEARCH METHOD

The development of Multiple Error Correction (MEC) consists of several methods. First, we proposed the MEC algorithms based on the Hybrid ARQ (HARQ) process. Second, we aided the algorithm by calibrating the transmit power according to the MEC assigned for different error correcting codes at different SNR range. Lastly, we integrated Kalman Filter (KF) to estimate the SNR values for incoming transmission in order to avoid wrong assignment of error correcting codes towards the sudden changing environment of CDMA WSN.

Figure 1 shows the chronological development of our proposed algorithm. We also carried out three different preliminary experiments comprising the simulation of CDMA WSN with no error correction, with a range of low error correcting capability, and with a range of high error correcting capability. We selected BCH and RS codes as the error correcting codes in our algorithm based on the preliminary experiments where BCH and RS codes optimized the performance of CDMA WSN compared to the Convolutional Codes [11], [12]. Some studies also agreed that the RS [14] and BCH codes [15] are among other optimum error correcting codes which provide reliability during transmission. From our preliminary testing [11]-[13], we observed that in a low SNR condition with high BER, the use of low error correcting capability might not be able to solve the problem of high BER. Meanwhile, the use of high error correcting capability caused the problems of increment in energy consumption as well as latency. We also found that medium error correcting capability for the high SNR condition might be able to optimize BER performance as well as the remaining energy. Besides that, the use of codeword length corresponding to the number of generated bits in the network is also critical. Higher numbers of codeword lengths (more than 127) promote more energy usage.
From this observation, we concluded that there is the need to use multiple coding schemes for different range of SNR. We also considered some of other related works [16] which demonstrated that the multi-coding is more energy-efficient. Based on this deduction, we proposed the SNR classification algorithm in adapting to the changes of channel condition. This algorithm follows the existing process of HARQ in which we modified the algorithm to be injected with MEC module and transmit power module. We also integrated Kalman Filter (KF) to estimate the SNR values to avoid assigning wrong error correction due to sudden changes in a noisy channel. Some field tests [17]-[19] observed that the signal impairments is time-varying. In a time-varying condition, there is a need to continuously track the changes within the channel as signals frequently attenuated, reflected, and refracted. In addition, the hardware itself might be a contributing factor to the changes in SNR as different nodes gives out different SNR. Thus, channel estimation is substantial to ensure that proper error correcting codes assigned to a defined SNR range which corresponds to the sudden changes in SNR when there are transmissions between nodes in any given time. KF is among the common state estimator with advantages such that low complexity compared to other filters [20] and its capability to provide better estimation for Gaussian and linear models, and limited non-linearity. We integrated KF prior to generating data and after the decoding process. In order to minimize the computational overhead of KF equation, we modified the ACK/NACK message which append the SNR value of the estimation initiated to the transmitted bits from the sender. This is to acknowledge the receiver on the previous estimation value and to prepare the transmission if there are sudden changes in SNR value so that the possibility in assigning wrong error correcting codes is reduced if the sender initiates another transmission.

2.1. SNR Classification towards the Adaptation to the Changes in Channel Condition

We divided the methodology into two following solutions based on the unsolved problems of high BER and the problems from high energy usage and high latency due to the implementation of high error
correcting capability (Figure 1). In method 1, we proposed the channel adaptation algorithm based on HARQ process and classified the range of SNR. We calculated the range of SNR for our predefined topologies of random uniform node distribution and random non-uniform node distribution. A random uniform node distribution consists of layered topology with random node scattering around the sink with uniform distance of 10 m between each node. Meanwhile, the random non-uniform distribution consists of layered topologies where the distance between nodes varied between 10–100 m. The SNR was calculated following the standardised equation below [21]:

\[
\text{SNR (dB)} = P_{\text{signal}} - P_{\text{noise}}
\]

where \(P_{\text{signal}}\) denotes the received signal power while \(P_{\text{noise}}\) denotes the noise power in which that \(P\) is measured in units of power (Watts or milliwatts). Extending the formula, the signal power can be obtained using Friis Transmission equation as shown in Equation (2) [22]:

\[
P_{\text{signal}} = \frac{(P_T G_T G_R) (\lambda^2)}{(4\pi R)^2}
\]

where \(P_T\) is transmit power, \(G_T\) is the gain of transmit antenna, \(G_R\) is the gain of receiving antenna, and \(R\) is the distance between sender and receiver. The noise power, \(P_{\text{noise}}\) is calculated following Equation (3) [23].

\[
N_0 = k T_s B
\]

where \(k\) denotes the Boltzmann constant, \(T_s\) denotes the system temperature, and \(B\) denotes the bandwidth.

From the calculation, we obtained the actual SNR that follows our predefined architecture and parameters as shown in Table 1 (Section 2.3). We extracted and applied the calculated values into the pre-built system identification model using MATLAB. From the extracted data, we obtained the matrices from the best fit model using the input-output model from the system identification tool; matrices \(A\) (-0.236468303496098), \(B\) (0.113429553384058), and \(C\) (3.59450124244613). Table 1 shows the extracted information of actual SNR alongside with the estimated SNR obtained using Kalman Filter equation as shown in Equation (4) and Equation (5).

\[
X_k = A x_{k-1} + B u_{k-1} + w_{k-1}
\]

\[
Y_k = H x_k + v_k
\]

where \(A\) is the state transition model applied to the previous state \(x_{k-1}\), \(B\) is the control-input model applied to the control vector \(u_{k-1}\), and \(w\) is the process noise. \(H x_k\) denotes the observation model which maps the true state space into the observed space and \(v_k\).

Table 1 shows the actual SNR obtained with the estimated SNR values after we have calculated the measurement and innovative gain from KF equation. We then divided the SNR range into five distinct groups: \(\text{SNR}_{\text{low}}\) (SNR lowest) with the range of \(\text{SNR}_{\text{v}} \leq 5\), \(\text{SNR}_{\text{low}}\) (SNR low) with \(6 \leq \text{SNR}_{\text{v}} \leq 20\), \(\text{SNR}_{\text{medium}}\) (SNR medium) with \(21 \leq \text{SNR}_{\text{v}} \leq 35\), \(\text{SNR}_{\text{high}}\) (SNR High) with \(36 \leq \text{SNR}_{\text{v}} \leq 50\), and \(\text{SNR}_{\text{high}}\) (SNR highest) where \(\text{SNR}_{\text{v}} \geq 51\). Based on previous studies and real-time testing [24], most studies agreed that the link is in a good quality when the SNR value is more than 40 dB. Thus, we set the benchmark of SNR in high condition (SNR\(_\text{h}\)) which is considered as good link quality. According to the study [25], SNR value below 5 dB is considered as bad link quality as the Packet Delivery Ratio (PDR) is zero. According to some studies, SNR alone is not sufficient to indicate congestion when high number of nodes present in one monitoring area. Thus, we implemented the congestion detection based on increasing node density as a network will be congested when node density increases above certain point. A network is assumed to be congested when the range of node density is more than 32 nodes in one given area based on a previous study [26]. The PDR and throughputs started to decrease indicating congestion as the node density reaches 45. Therefore, we noticed that the network started to congest when it reached 48 nodes for our pre-defined topology (300 m x 300 m) of the monitoring area. We extended the classes from previous SNR range mentioned where we indicated that the network is extremely congested (EC) when the node density is more than 27.9253, medium congested (MC) when the range of node density is between 16.7552 and 22.3402, and non-congested (NC) when the node density is less than 11.1701.
2.2. Selection of the Optimal Error Correcting Codes and Transmit Power

Method 2 was divided into three sub-methods of Method 2.1, Method 2.2, and Method 2.3. Following the classification of SNR range in Method 1, we simulated the error correction codes of BCH and RS with a variety of error correcting capabilities such that $1 \leq t \leq 10$. We implemented two types of error correcting codes (BCH and RS codes) following the range of codeword length, $n$ denoted as $15 \leq n \leq 127$ for random uniform distribution and $63 \leq n \leq 127$ for random non-uniform node distribution. The idea is to simulate each error correcting code in obtaining the most optimal code with corresponding codeword length and error correcting capability in order to adapt to different channel conditions. For instance, our proposed algorithm implemented high error correcting capability for lower SNR to significantly reduce the too high BER whereas for a good network condition, we implemented a much lower error correcting capability. The reason is that when the SNR condition is good (good BER performance), too high error correcting capability might not be necessary as retransmission is enough to solve the errors. This reduces the overhead of encoding and decoding processes of the ECC as well as the latency caused by the computation of that ECC.

Table 2 and Table 3 show the lists of tested BCH and RS with respective codeword length and error correcting capability. The codes were simulated with CDMA WSN to study the effects of BER and remaining energy as well as latency and to find the most optimal error correcting codes for every range of SNR that had been classified. We also added the transmit power calibration to our proposed algorithm to maximize remaining energy and reduce noises in congested area (Method 2.2). In a condensed network where the nodes are so close with each other and having a good range of transmission coverage, implementing high transmit power might not be necessary.

In addition, in this condensed network, the nodes might be interfering with the transmission range of another nodes [8],[9]. This promotes the increase in interferences and noises and yet the reduction of transmit power might be able to reduce the Multiple Access Interference (MAI) from the CDMA architecture. In a condensed network where generally the high number of nodes with high transmit power might lead the MAI to peak [27],[28], causing the packet to be corrupted and increase the BER. The reduction in transmit power might be able to slightly increase the remaining energy [29] and theoretically might be able to reduce MAI as

| Distance (m) | Transmit Power (dBm) | Actual SNR (dB) | Estimated SNR (dB) |
|-------------|----------------------|----------------|-------------------|
| 10          | 0                    | 41.9764        | 41.9843           |
| 20          | 0                    | 31.4403        | 31.4482           |
| 30          | 0                    | 25.2771        | 25.2850           |
| 40          | 0                    | 20.9043        | 20.9122           |
| 50          | 0                    | 17.5124        | 17.5203           |
| 60          | 0                    | 14.7411        | 14.7490           |
| 70          | 0                    | 12.3979        | 12.4058           |
| 80          | 0                    | 10.3682        | 10.3761           |
| 85          | 0                    | 9.4467         | 9.4546            |
| 90          | 0                    | 8.5779         | 8.5858            |
| 100         | 0                    | 6.9764         | 6.9843            |

| Error Correcting Codes | $n$ | $k$ | Error Correcting Capabilities, $t$ |
|------------------------|-----|-----|-----------------------------------|
| BCH                    | 7   | 4   | 1                                 |
|                        | 15  | 5   | 3                                 |
|                        | 31  | 26  | 1                                 |
|                        | 31  | 16  | 3                                 |
|                        | 31  | 11  | 5                                 |
|                        | 63  | 57  | 1                                 |
|                        | 63  | 45  | 3                                 |
|                        | 63  | 24  | 7                                 |
|                        | 63  | 18  | 10                                |
|                        | 127 | 120 | 1                                 |
|                        | 127 | 92  | 5                                 |
|                        | 127 | 64  | 10                                |
|                        | 255 | 247 | 1                                 |
|                        | 255 | 179 | 10                                |
|                        | 511 | 502 | 1                                 |
|                        | 511 | 466 | 5                                 |
|                        | 511 | 421 | 10                                |

| Error Correcting Codes | $n$ | $k$ | Error Correcting Capabilities, $t$ |
|------------------------|-----|-----|-----------------------------------|
| RS                     | 7   | 3   | 1                                 |
|                        | 15  | 11  | 2                                 |
|                        | 15  | 5   | 5                                 |
|                        | 31  | 27  | 2                                 |
|                        | 31  | 13  | 9                                 |
|                        | 31  | 11  | 10                                |
|                        | 63  | 59  | 2                                 |
|                        | 63  | 45  | 9                                 |
|                        | 63  | 43  | 10                                |
|                        | 127 | 123 | 2                                 |
|                        | 127 | 113 | 7                                 |
|                        | 127 | 107 | 10                                |
|                        | 255 | 251 | 2                                 |
|                        | 255 | 235 | 10                                |
|                        | 511 | 507 | 2                                 |
|                        | 511 | 499 | 6                                 |
|                        | 511 | 491 | 10                                |
Multiple error correction towards optimisation of energy in sensor network (Samirah Razali)

well. However, we also suggested that the transmit power reduction must be optimal as too much reduction of transmit power will also causes packet drop as the signal power decreases whereas the distance between nodes in the non-random uniform distribution increase to 100 m. The transmit power can be calculated by using the Friis Transmission equation shown in Equation (2). Equation (6) denotes the simplified equation to obtain transmit power from the received signal power:

\[ P_t = \frac{P_r \times (4\pi d^2)}{d^2} \]  

(6)

where \( P_r \) is the received signal and \( d \) is the distance between nodes.

2.3. Measurement Models

Table 4 shows the parameters defined for testing and simulation of the proposed work. We implemented the minimum distance of 10 m for uniform distribution and 100 m for non-uniform distribution. Non-uniform distribution is referred to as the distance between nodes (might not be the same from one node to another) resulting in different levels of node density present within one monitoring area.

| Parameter                          | Value                                      |
|------------------------------------|--------------------------------------------|
| Min. dist. between two nodes       | 10 m                                       |
| Noise                              | AWGN                                       |
| Transmit Power (\( P_t \))         | Power allocation depending on SNR estimation (-10 dBm to 0 dBm) |
| Monitoring area (meter\(^2\))      | 300 m x 300 m                              |
| Path loss parameter (\( \alpha \)) | 3.5                                        |
| Payload, Header                    | 128,256 bit/pkt                            |
| NACK/ACK (\( n_b \))               | 8 bits (additional 6 bits as NACK/ACK were appended with SNR and update value from the receiver) |
| Error Detection                    | CRC-30 (CDMA compliance)                   |
| Error Correction                   | Proposed MEC (variation on the BCH and RS ECC) |
| Number of Nodes                    | 4,16,32,48,64,80                           |
| Number of Bits (bits)              | 10000, 20000, 30000, 40000, 50000          |

The application of 100 m (as the maximum distance between nodes) supports the common sensor nodes such as MicaZ and TelosB with the transmission range up to 100 m. We added Additive White Gaussian Noise (AWGN) and Rayleigh Fading to study the impacts between these two interventions. The power allocation was calibrated according to the SNR range as well as error correction schemes.

We developed our measurement model to measure the performances of CDMA WSN in terms of remaining energy, BER, and latency. The following equations were used to obtain the measurement results. The expression of remaining energy including decoding energy can be denoted as:

\[ \text{Energy}_{\text{Ecc}} = \text{Hop} \times \text{No_packet} \times (\text{No_bits} + (\text{No_bits} \times 0.75) + \text{Energy}_{\text{Dec}}) \]  

(7)

where Hop corresponds to number of hops of the predefined layered architecture, No_packet is the number of packets involved in the transmission, No_bits is number of generated bits which for research (10000 bits), and \( \text{Energy}_{\text{Dec}} \) is the decoding energy calculated in as shown in equation (8).

\[ \text{Energy}_{\text{Dec}} = (2n + 2t^2)(\text{Energy}_{\text{addition}} + \text{Energy}_{\text{multiplication}}) \]  

(8)

where \( n \) denotes the codeword length of the selected error correcting codes and \( t \) is the error correcting capability of that error correcting codes. We tabulated the tested codes in which were used in the calculation of decoding energy in Table 2 and Table 3. We also modelled the BER formulation in BPSK modulation as follows:

\[ \text{BER} = \frac{1}{2} \left(1 - \sqrt{\frac{E_b}{N_0}} + \frac{E_b}{N_0} \right) \]  

(9)

Latency is calculated following Equation (10), where \( \text{Delay}_{\text{propagation}} \) is the propagation delay. Propagation delay refers to the time taken between departure of data from the sender and arrival of data at the
receiver. Meanwhile, $Delay_{Transmission}$ is the transmission delay or also known as the packetization delay in which can be defined as the amount of time required to transmit all of the packet's bits into the link. Usually, the transmission delay is affected by data-rate of the link. Propagation delay was calculated by dividing the distance between sender and receiver with propagation speed of the media. The transmission delay was calculated by dividing the length of packet in bits with the transmission rate of the predefined network.

$$Latency = Delay_{Propagation} + Delay_{Transmission}$$ (10)

2.4. Optimal Error Correction Schemes for the proposed Multiple Error Correction (MEC)

Table 5 shows the settings of optimal error correction schemes for the proposed Multiple Error Correction (MEC) algorithm. This setting is optimal if the architecture follows our predefined topologies in which the default number of bits is 10000 bits with increasing number of nodes (between 4 and 80 nodes). In this study, the CDMA architecture implemented Multi-Carrier (MC-CDMA) channel access method and BPSK modulation as tabulated in Table 4. From the simulation results acquired, the BCH codes with codeword length of 127 cannot outperform RS codes with the same codeword length when the network condition is extremely congested due to redundancy added by the BCH codes is higher than the RS codes. In some studies, more redundancies might lower the BER. However, in our case, too high redundancy appended to the transmitted bits in a congested network flooded the network causing more bits to be corrupted. Thus, we observed that RS code with codeword length of 127 is optimum for an extremely congested condition. For a medium congested condition, we applied BCH codes for both, high SNR and the highest SNR with good link condition where BER is relatively low. Higher appended bits might not corrupt the bits as much as when the SNR is low. For a non-congested condition, we applied retransmission for higher SNR with low BER present in the network as the retransmission is enough to correct the erroneous bits. This is also to add up that there were only a few nodes that might interfere with the transmission range in which if there were collisions or interference that can corrupt the bits, the errors might not be too significant compared to the condition with higher node density.

| Condition: EC | Error Correcting Codes | Codeword length, $n$ | Information Bits, $k$ | Error Correcting Capability, $t$ |
|---------------|-------------------------|----------------------|-----------------------|-------------------------------|
| Case 1: ECSNR_lw | RS | 127 | 113 | 7 |
| Case 2: ECSNR_l | RS | 127 | 115 | 6 |
| Case 3: ECSNR_m | RS | 127 | 117 | 6 |
| Case 4: ECSNR_h | RS | 127 | 119 | 4 |
| Case 5: ECSNR_hg | RS | 127 | 121 | 3 |

| Condition: MC | Error Correcting Codes | Codeword length, $n$ | Information Bits, $k$ | Error Correcting Capability, $t$ |
|---------------|-------------------------|----------------------|-----------------------|-------------------------------|
| Case 1: MCSNR_lw | RS | 63 | 53 | 5 |
| Case 2: MCSNR_l | RS | 63 | 51 | 6 |
| Case 3: MCSNR_m | RS | 63 | 49 | 7 |
| Case 4: MCSNR_h | BCH | 63 | 36 | 5 |
| Case 5: MCSNR_hg | BCH | 63 | 24 | 7 |

| Condition: NC | Error Correcting Codes | Codeword length, $n$ | Information Bits, $k$ | Error Correcting Capability, $t$ |
|---------------|-------------------------|----------------------|-----------------------|-------------------------------|
| Case 1: NCSNR_lw | BCH | 127 | 64 | 10 |
| Case 2: NCSNR_l | RS | 127 | 121 | 3 |
| Case 3: NCSNR_m | BCH | 31 | 11 | 6 |
| Case 4: NCSNR_h | RS | 31 | 17 | 7 |
| Case 5: NCSNR_hg | Retransmission with IR |

Table 6 shows the optimal transmit power corresponding to different classes and congestion present based in our abovementioned architecture. The transmit power with high SNR were significantly reduced to 20% for extremely congested condition, 40% for medium congested, and 50% for non-congested. In the extremely congested condition, the transmit power was not reduced too much as compared to the non-congested condition in order to avoid increment in the noise present. This is because the signal power will reduce the transmit power calibrated. The lower power will be peak the noise, which lead to the increment in BER. For non-congested condition, the transmit power can be reduced until 50% as having high SNR indicate high signal power with low noise power.
Table 6. Optimal transmit power for MEC

| Condition: EC | Transmit power calibration |
|---------------|---------------------------|
| Case 1: ECSNR_{lw} | Maintain |
| Case 2: ECSNR_{l} | Maintain |
| Case 3: ECSNR_{m} | Reduction by 10% |
| Case 4: ECSNR_{h} | Reduction by 10% |
| Case 5: ECSNR_{hg} | Reduction by 20% |

| Condition: MC | Transmit power calibration |
|---------------|---------------------------|
| Case 1: MCSNR_{lw} | Maintain |
| Case 2: MCSNR_{l} | Reduction by 10% |
| Case 3: MCSNR_{m} | Reduction by 20% |
| Case 4: MCSNR_{h} | Reduction by 30% |
| Case 5: MCSNR_{hg} | Reduction by 40% |

| Condition: NC | Transmit power calibration |
|---------------|---------------------------|
| Case 1: NCSNR_{lw} | Reduction by 10% |
| Case 2: NCSNR_{l} | Reduction by 20% |
| Case 3: NCSNR_{m} | Reduction by 30% |
| Case 4: NCSNR_{h} | Reduction by 40% |
| Case 5: NCSNR_{hg} | Reduction by 50% |

3. RESULTS AND ANALYSIS

We collected the results based on the experiment of MEC in terms of increment in node density and compared the non-channel adaptation using the existing error correction schemes. Figure 2 shows the average remaining energy against SNR between MEC with BCH (127, k) t = 10, RS (127, 121) t = 3, BCH (31, k) t = 7, and RS (31, k) t = 7. The existing RS (31, k) t = 7 have the highest constant remaining energy throughout the SNR. The highest remaining energy for MEC is obtained when the SNR is at the good condition. Even though RS (31, k) t = 7 had the highest remaining energy in a non-congested condition, the higher redundancy added by this code was not practical for higher node density.

Figure 2. Average Remaining Energy against SNR

In addition, MEC cannot reduce BER as further reduction of the remaining energy was too much due to drawbacks between the capabilities of error corrections that added redundancies with more complexities and required higher energy when better error correction is used. However, MEC optimized the remaining energy in high SNR as too powerful error correction schemes were not necessary. Instead, the retransmission was enough to handle the errors which do not used too much energy as compared to ECC due to decoding and computation overhead. However, for low SNR (where BER is relatively very high), it is substantial to correct the errors to reduce too many resend packets from the sender. A higher number of
flooded retransmissions in the network not only uses energy but also resources as well which might render the network causing malfunction. The remaining energy for low SNR was lesser compared to other SNR range as higher error correction capability was used. Thus, we can conclude that the optimisation between remaining energy and BER by means of adaptation to the channel condition itself. In low SNR, the remaining energy might not be able to be boosted as much as when the SNR is high. This low error correcting capability codes might not able to handle the high error rates if lower error correcting capability is used in low SNR where BER is high.

Figure 3 shows the comparison of average latency against SNR between MEC and existing codes. The existing coding schemes have constant latency throughout the SNR due to the absence of channel adaptation. The latency changed with SNR range due to the changes in error correction used throughout the SNR as shown in the figure. The latency corresponded to the appended bits in the network where higher appended bits or redundancy causes higher latency as large number of bits prolonged the decoding time of the received information. In this study, BCH recorded higher latency due to its complexity and higher number of redundancies. However, this corresponds to the fact that BCH have better error correction capability that the RS code. For instance, BCH (31, k) outperformed RS (31,k) with the same t = 7 in terms of BER even though the latency was too high. However, the use of multi-coding might able to optimize the latency as the latency will not always be high throughout the SNR. When high redundancy is not needed, the latency will be reduced by using much lower capability codes that appended much lower redundancy.

Figure 3. Average Latency against SNR

Figure 4 illustrated the average BER against increasing node density. The MEC showed the lowest BER compared to existing schemes for the non-congested condition of less than 11.1701 node density when SNR reaching 30 dB. The reason is that BCH (31, k) was implemented for medium SNR in non-congested environment as such redundancy significantly solve the BER and does not flood the network in the non-congested environment. However, for higher node density, BCH (31, k) was impractical because of too high redundancy. However, the existing RS (127, k) did not outperform BER of MEC for non-congested condition (≤ 11.1701).

Figure 5 shows the average remaining energy against node density. The network achieved the highest remaining energy in congested condition for medium SNR as we used medium error correcting capability which optimised the need between remaining energy and redundancy added to the network. There was much higher error correcting capability used for low congested condition as the remaining energy was slightly lower than the existing RS (31, k). However, in comparison to Figure 4 where BER of MEC was at its lowest when node density at 11.1701, the higher remaining energy of RS proved that the existing RS did not optimize the BER and remaining energy. Meanwhile, there was slight increment in the remaining energy from 628946.33 J for node density between 16.7552 and 22.3402.
Figure 4. Average BER against Node density (for SNR=30)

Figure 5. Average Remaining Energy against Node density

Figure 6 shows the Percentage increment in remaining energy of MEC against SNR compared with BCH (127,k) for t=7, RS (127,k) for the same t and the comparison of increment in remaining energy between existing RS and BCH. From the graph, it can be observed that as the SNR increased, the percentage increment in remaining energy also increased. Besides that, the percentage of increment of remaining energy of MEC is higher than the existing RS (127, k) and BCH (127, k) when SNR achieving between 20 dB and 50 dB.

It is also pointed out that the RS (127, k) has a higher percentage of increment during the low SNR of 20 dB because MEC has more powerful error correction schemes which used more energy. However, the difference between the percentage increment of MEC and RS is quite small. This also corresponds to the graph in Figure 4 which demonstrated that during the lowest SNR, MEC achieved better BER than the existing RS which results in lower remaining energy. Thus, a conclusion was made where MEC had optimised the performances between the BER and remaining energy and for the lowest SNR and for highest SNR, MEC performed more towards optimisation of remaining energy and latency.
4. CONCLUSION

This paper developed the algorithm towards adapting the changes in channel condition and congestion by classifying SNR into several ranges in order to optimise the performances between remaining energy and BER. Our proposed algorithm merits on its capability to adapt to the changes in SNR and optimise the performance of remaining energy and BER in low SNR as well as high SNR in which every SNR range describe different error correction used and transmit power calibration to minimise the higher redundancies when these excessive redundancies are not needed in better SNR condition. Transmit power calibration aided to reduce the interferences and noises in noisy channel when SNR is low. For example, when SNR was at its highest, we assumed the network condition was good and retransmission was likely more sufficient to correct the corrupted bits compared to error correcting codes. While, when SNR is low, the noise ratio usually fluctuated and causes the BER to rises which in need for higher error correcting capability. Relating to our analysis, we had obtained lower BER during low SNR compared to the existing method in which the error correction was constant throughout the increasing SNR. While, we are able to maintain remaining energy and BER by using medium error correcting capability during the medium and high SNR condition as for existing method, higher error correcting capability such as BCH (127,k) might not be necessary as demonstrated have the lowest remaining energy in Figure 2 even when the SNR is high. In addition, MEC outperformed the existing BCH as the number of nodes increased because MEC implemented RS (63, 57) during medium SNR corresponding to the fact that the RS codes in MEC possessed lower latency than BCH due to its lower error correction capability that appended much lower number of redundancies than BCH. However, in a medium network condition of SNR, it was observed that RS corrected the errors as well as BCH since the BER was optimised and significantly reduced.

As for the future work, it is suggested to study the multi-coding schemes of MEC with other error correcting codes such as Turbo codes, Hamming, and Low-Density Parity Check (LDPC) codes. Last but not least, it is recommended that future research to include the aspects of localisation of nodes as well as mobility since the study of localisation promotes better scalability with different geographical limitations assessment. Besides, mobility might further improve MEC in supporting different types of environmental monitoring other than natural disasters such as pollution monitoring, agricultural monitoring, and even human sensing that requires node mobility features.

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