Low power pseudo-random number generator based on lemniscate chaotic map

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ABSTRACT

Lemniscate chaotic map (LCM) provides a wide range of control parameters, canceling the need for several rounds of substitutions, and excellent performance in the confusion process. Unfortunately, the hardware model of LCM is complex and consumes high power. This paper presents a proposed low power hardware model of LCM called practical lemniscate chaotic map (P-LCM) depending on trigonometric identities to reduce the complexity of the conventional model. The hardware model designed and implemented into the field programmable gate array (FPGA) board, Spartan-6 SLX45FGG484-3. The proposed model achieves a 48.3% reduction in used resources and a 34.6% reduction in power consumption compared to the conventional LCM. We also introduce a new pseudo-random number generator based on a proposed low power P-LCM model and perform the randomization tests for the proposed encryption system.

Keywords: Chaotic map, FPGA, Lemniscate chaotic map, Random number generator, Read-only memory (ROM)

INTRODUCTION

Pseudo chaotic random number generators (PCRNG) is a vital cryptography application of nonlinear chaotic systems. Many researchers present the software implementation of PCRNGs using MATLAB. The most common way to generate PCRNG is to use a feedback shift register method in different ways such as a linear feedback shift register, carry forward feedback shift register, and nonlinear feedback shift register [1-3]. Another technique used a coupled map lattice with time-varying delay [4].

The researchers in [5] proposed a PRNG based on Lorenz systems with FPGA implementation achieving an operating frequency of 78.149 MHz. In [6], the authors investigate the fixed-point arithmetic representation random effect, and they presented a PRNG based on coupled skew tent maps and provides its hardware implementation. The FPGA implementations of different PRNGs based on chaotic maps presented in [7]. In [8], A pseudo-random generator generated using a chaotic quadratic map with FPGA implementation. In [9], another idea used to make PRNG with a function based on ring oscillator and chaotic logistic map. Another design to PRNG based on a logistic map that changes its chaotic parameters presented and implemented with FPGA in [10]. In [11], a secured PRNG based on a piecewise linear chaotic map presented and implemented using FPGA.

The sensitivity of the chaos-based encryption systems is affected by two factors; the first factor is the use of binary streams extracted from a single orbit of a chaotic map while the second factor is the use of

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maps that have chaotic behavior only for small ranges of control parameters’ values. Another issue that affects the chaos-based encryption systems is a low speed, which can be caused by the need for several rounds of permutation and/or substitution of the original image pixel [12]. The LCM whose cryptographic properties have been demonstrated to be very good in the confusion process, and eliminate the need for several rounds of substitutions of the pixels values and have a wide range of control parameter’s values [13]. The basic equations of LCM are:

\[ x(n + 1) = \frac{\cos(2^r y(n))}{1 + \sin^2(2^r y(n))} \]  \hspace{1cm} (1)

\[ y(n + 1) = \frac{2\sqrt{2} \sin(2^r x(n)) \cos(2^r)x(n)}{1 + \sin^2(2^r x(n))} \]  \hspace{1cm} (2)

Where the initial conditions \( x_0, y_0 \) are given from the interval \([-1: +1]\) and \( r_0 > 3 \) for the hyperchaotic regime. In Figure 1, the bifurcation diagram and the Lyapunov exponents of LCM are shown in Figure 1 (a) and Figure 1(b) respectively.

![Figure 1](image)

(a) Bifurcation diagram, (b) Lyapunov exponents

The equations of LCM have trigonometric functions (sin, cos), and it can be implemented in a hardware model by different methods such as: look-up table (LUT) based read-only memory (ROM), CORDIC algorithm, Taylor’s series, and linear segmentation [14-16]. Since the total power consumption of a hardware model depends on the components used in the implementation, so any reduction for these components reduces the overall power consumption. In the LCM hardware model, the components LUT, ROM, consume more power so, the total power consumption reduction achieved by reducing the number of these components in the hardware model. The main idea of this paper is to present a low power implementation of the lemniscate hardware model. Three alternatives hardware models designed, explained, and implemented according to the mathematical equations of LCM. Furthermore, we present a PRNG based on the best low power FPGA architecture of the LCM with MATLAB simulations and FPGA simulation and implementation.

This paper organized as follows: Section 2 describes the mathematical analysis of three alternative models of LCM. In section 3, FPGA Implementations of the three hardware models of LCM are presented. The FPGA hardware implementation results provided in section 4. The hardware implementation of the PRNG based on the lowest power consumption hardware model is presented in section 5. Section 6 presents NIST SP800-22 randomization tests for a proposed PRNG. Section 7 presents a comparison between the proposed model and recent comparable models. Finally, this conclusion presented in section 8.
2. PROPOSED LCM MODELS

2.1. Modified model

The first alternative “Conventional” model achieved by using the basic mathematical equations of lemniscate chaotic map, which are (1-2). The second alternative “Modified” map uses two abbreviations to (1) and (2) using the trigonometric identities, which are:

\[
\sin^2(t) + \cos^2(t) = 1
\]

\[
2 \sin(t) \cos(t) = \sin(2t)
\]

Using (3) converts the denominator of (1) as follows:

\[
1 + \sin^2(2^r y(n)) = 2 - \cos^2(2^r y(n))
\]

Also, using (4) converts the numerator of (2) as follows:

\[
2\sqrt{2} \sin(2^r x(n)) \cos(2^r x(n)) = \sqrt{2} \sin(2 \times 2^r x(n))
\]

So, the second alternative “Modified” has the following equations

\[
x(n+1) = \frac{\cos(2^r y(n))}{2 - \cos^2(2^r y(n))}
\]

\[
y(n+1) = \frac{\sqrt{2} \sin(2 \times 2^r x(n))}{1 + \sin^2(2^r x(n))}
\]

2.2. Practical model

The third alternative “Practical” uses three abbreviations to (1) and (2); using the trigonometric identities in (3), (4), and the following identity:

\[
\cos(2t) = \cos^2(t) - \sin^2(t) = 2 \cos^2(t) - 1 = 1 - 2 \sin^2(t)
\]

\[
\sin^2(t) = \frac{1}{2} - \frac{1}{2} \cos(2t)
\]

In (3) converts (1) to be (5), while (4) converts (2) to (6), as done in “Modified” model. Substitute with (7) in (6) leads to

\[
y(n+1) = \frac{\sqrt{2} \sin(2 \times 2^r x(n))}{1.5 - 0.5 \cos(2 \times 2^r x(n))}
\]

Since

\[
\frac{d}{dx(n)}(\sin(2 \times 2^r x(n)) = 2 \times 2^r \cos(2 \times 2^r x(n))
\]

Rewriting (8)

\[
y(n+1) = \frac{\sqrt{2} \sin(2 \times 2^r x(n))}{1 + \sin^2(2^r x(n))}
\]

In discrete-time, the time differentiation is calculated using

\[
\frac{d}{dn} x(n) = x(n) - x(n-1)
\]
It means by only using a one-time delay and subtractor, we can obtain a time differentiator. So, the third alternative “practical” hardware model implements the following equations:

\[
x(n + 1) = \frac{\cos(2^r y(n))}{2 - \cos^2(2^r y(n))}
\]

(15)

\[
y(n + 1) = \frac{\sqrt{2} \sin(2 \times 2^r x(n))}{1.5 - 2^r \times \frac{d}{dx(n)}(\sin(2 \times 2^r x(n)))}
\]

(16)

3. FPGA IMPLEMENTATION OF PROPOSED LCM

The hardware models of the lemniscate architectures is designed using Xilinx system generator (XSG) [17-22]. In this paper, we implement the sine and cosine functions using LUT based ROM in the three hardware models. Table 1 indicates the number of LUT used in each alternative.

| Model      | No. of LUTs |
|------------|-------------|
| Conventional | 4           |
| Modified    | 3           |
| Practical   | 2           |

The three alternatives of the LCM system are modelled using the XSG program in thirty two fixed-point formats, also the three architectures are implemented into the same FPGA board (Spartan-6 SLX45FGG484-3). The XSG conventional hardware architecture model which depend on (1-2) is shown in Figure 2. The “Modified” XSG model, which implement (7-8) is shown in Figure 3. The “Practical” XSG model which applies the (15-16) is shown in Figure 4.

![Figure 2. XSG model of “Conventional” LCM](image-url)
4. HARDWARE SIMULATION RESULTS

4.1. Bifurcation diagram

Table 2 shows a comparison between the bifurcation diagram for the three models generated by theoretical simulations using Matlab and bifurcation diagram generated by the hardware model. The difference between the two types of simulation because of the finite word length in the digital hardware models [23].

4.2. Implementation results

The implementation results for the three-hardware LCM models are presented in Table 3. Note that the maximum frequency is the same for the three hardware models because the maximum frequency is determined by the critical path, which is the longest path between the input and output signal. Since the steps of calculations are the same in the three hardware models, so the maximum frequency is the same.
Table 2. Comparison between Bifurcation diagrams generated by Matlab codes, and by hardware models

| Model      | Theoretical Simulation (Matlab Code) | Hardware Model Simulation (VHDL code) |
|------------|-------------------------------------|---------------------------------------|
| Conventional | ![Diagram](image1.png) | ![Diagram](image2.png) |
| Modified   | ![Diagram](image3.png) | ![Diagram](image4.png) |
| Practical  | ![Diagram](image5.png) | ![Diagram](image6.png) |

Table 3. Hardware implementation results of LCM

| Model      | Slice Registers | LUTs | Frequency (MHz) | Power (mW) | Power reduction ratio Compared to Conventional |
|------------|-----------------|------|-----------------|------------|-----------------------------------------------|
| Conventional | 1247            | 7614 | 80.592          | 367        | ---                                           |
| Modified   | 1024            | 5806 | 300             | 300        | 18%                                           |
| Practical  | 859             | 3329 | 240             | 240        | 34.6%                                         |

5. FPGA IMPLEMENTATION OF P-LCM RNG

An XSG hardware model of the random number generator is built by adding two threshold units and a multiplexer to the outputs of the “Practical” model (x,y) as shown in Figure 5. The threshold is a “Relational” block compares the input to a threshold value which is “0.5”, if the input is a fraction higher than “0.5” then the output of the “Relational” block is “1”, otherwise the output will be “0”. Figure 6 shows a simulation for RNG. A random output number is a binary number “0” or “1” according to the “sel” signal which chooses between the inputs “d0” and “d1”.

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6. NIST SP800-200 RANDOMIZATION TESTS

A good encryption system should be able to assign plane images to randomly encoded images. It is, therefore, vital to test the randomization of the encrypted images obtained by the proposed image encryption algorithm [24]. The results given in Table 4 indicates that the P-LCM PRNG is suitable for cryptographic applications.

| Statistic Tests                   | p value  | Results |
|----------------------------------|----------|---------|
| Runs test.                       | 0.2219   | Pass    |
| Longest run of ones             | 0.6638   | Pass    |
| Binary matrix rank test         | 0.7214   | Pass    |
| FFT (spectral)                   | 0.9293   | Pass    |
| Non-overlapping template        | 0.5495   | Pass    |
| Overlapping templates           | 0.6257   | Pass    |
| Universal statistical           | 0.6127   | Pass    |
| Linear complexity               | 0.8895   | Pass    |
| Serial test (1)                 | 0.9428   | Pass    |
| Serial test (2)                 | 0.7215   | Pass    |
| Approximate entropy             | 0.7621   | Pass    |
| Cumulative sums (Forward)       | 0.8164   | Pass    |
| Cumulative sums (Reverse)       | 0.6845   | Pass    |
| Random excursions               | 0.6732   | Pass    |
| Random excursions variant       | 0.2765   | Pass    |
7. COMPARISONS

In this section, the proposed encryption system compared to recent chaotic systems. Table 5 presents the security test results in the case of using the Lena image. As indicated in Table 5, the results of our encryption system are close to the recent systems. Another comparison between the three hardware models of LCM and other similar works based on chaotic maps are presented in Table 6.

| References | Entropy | Unified averaged changed intensity | Correlation | Number of changing pixel rate |
|------------|---------|-----------------------------------|-------------|-----------------------------|
| Z. Hau [25] | 7.9956  | 33.418                            | 0.0209      | 99.630                      |
| Z. Tang [26] | 7.990   | 33.390                            | 0.0857      | 99.600                      |
| Z. Deng [27] | 7.9931  | 33.365                            | 0.0032      | 99.5995                     |
| Our System  | 7.9980  | 33.448                            | 0.0014      | 99.661                      |

8. CONCLUSION

A solution to the high-power consumption problem of the LCM hardware model is explained, analyzed, and implemented. Using trigonometric identities two alternatives hardware models; “Modified”, and “Practical” have been presented instead of the “Conventional” model. The implementation results indicate 18% reduction in power consumption in case of using the “Modified” model instead of “conventional” model. Also, implementation results indicate a 33% reduction in power consumption when using the “Practical” model instead of “conventional” model. We proposed a new Pseudo number generator based on a proposed “Practical” model. Also, statistical analysis has been used to determine the randomization tests. Finally, two comparisons are presented; the first between the security of the proposed encryption system and similar recent chaotic cryptosystems, while the second between the hardware implementation models of the proposed LCM systems and similar recent works.

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