Smooth Bandit Optimization: Generalization to Hölder Space

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Abstract

We consider bandit optimization of a smooth reward function, where the goal is cumulative regret minimization. This problem has been studied for $\alpha$-Hölder continuous (including Lipschitz) functions with $0 < \alpha \leq 1$. Our main result is in generalization of the reward function to Hölder space with exponent $\alpha > 1$ to bridge the gap between Lipschitz bandits and infinitely-differentiable models such as linear bandits. For Hölder continuous functions, approaches based on random sampling in bins of a discretized domain suffices as optimal. In contrast, we propose a class of two-layer algorithms that deploy misspecified linear/polynomial bandit algorithms in bins. We demonstrate that the proposed algorithm can exploit higher-order smoothness of the function by deriving a regret upper bound of $\tilde{O}(T^{\frac{d+\alpha}{d+2}})$ for when $\alpha > 1$, which matches existing lower bound. We also study adaptation to unknown function smoothness over a continuous scale of Hölder spaces indexed by $\alpha$, with a bandit model selection approach applied with our proposed two-layer algorithms. We show that it achieves regret rate that matches the existing lower bound for adaptation within the $\alpha \leq 1$ subset.

1 Introduction

This paper considers the problem of black-box optimization of a reward function $f : \mathcal{X} \to \mathcal{R}$, that is bounded and defined on a compact $d$-dimensional domain $\mathcal{X}$, using active queries. At each round, the learner chooses an action $x_t$ by leveraging the previously collected data and observes a noisy and zeroth order feedback of the function value $f(x_t)$. In the bandit setting, the goal is to minimize the cumulative regret with respect to global maxima. This is also known as the continuum-armed bandit problem. The bandit framework is different from standard global zeroth order optimization because of its unique exploration-exploitation dilemma. While in zeroth order optimization problems, pure exploration will often suffice since the performance is measured by simple regret (i.e. difference between the optimized function value and true function maxima), in bandit optimization, the queried function values need to be controlled through the entire optimization process to minimize the cumulative regret. Therefore, the algorithms require different and often more careful design.

Most existing works on continuum-armed bandit optimization either assume parametric models such as linear bandits ([Dani et al. (2008); Abbasi-Yadkori et al. (2011); Rusmevichientong and Tsitsiklis (2010)]) for the reward function, or a black-box model where the reward function is assumed to be $\alpha$-Hölder continuous (including Lipschitz) with $0 < \alpha \leq 1$ with respect to some known metric ([Kleinberg (2005); Auer et al. (2007); Kleinberg et al. (2008); Bubeck et al. (2010); 2011; Locatelli and Carpentier (2018)]). The main purpose of this paper is to extend this assumption to the more general Hölder function space (definition[I]) with exponent $\alpha > 1$ and exploit the higher
order of function smoothness. Generalization to $\alpha > 1$ is a parallel to the Hölder assumptions in fundamental results in nonparametric regression (Stone (1982)), which has been used in a variety of applications such as economics (Yatchew (1998)). Approaches based on fitting an appropriate function using random samples in bins of a discretization of the domain (i.e., exploration) suffice as optimal for controlling cumulative regret for Hölder continuous reward functions with $\alpha \leq 1$, as well as controlling simple regret of Hölder smooth reward functions with any $\alpha > 0$. In contrast, controlling cumulative regret for Hölder smooth reward functions with $\alpha > 1$ requires finer control in bins over the queried values via a local exploration-exploitation tradeoff. Thus, instead of using a single layer algorithm that randomly samples from selected bins, we propose a class of algorithms that use two layers of bandit algorithms - one multi-armed bandit algorithm operating over the bins, and another set of misspecified linear/polynomial bandit algorithms operating in each bin to govern the local exploration-exploitation tradeoff. We derive regret bounds for this class of two-layer bandit algorithms and show that they match the existing lower bounds apart from log factors.

Additionally, we study the problem of adaptation to smoothness exponent $\alpha$ for a continuous scale of Hölder spaces. Unlike the simple regret minimization setting where this adaptation comes at no cost in terms of the minimax rates, it was shown by Locatelli and Carpentier (2018) that it is generally impossible to achieve minimax adaptation under cumulative regret. We propose a procedure with regret bound that matches the existing adaptive lower bound with only access to the range of the unknown parameter $\alpha$. We start by describing related works, followed by a summary of our contributions.

1.1 Related Works

Continuum-armed Bandit. In continuum-armed bandit problems, the domain $\mathcal{X}$ is allowed to be a measurable space, and the set of arms is therefore infinite. Previous works in continuum-armed bandit usually assume global smoothness (Kleinberg (2005)) of the reward function or local smoothness (e.g. Auer et al. (2007)) around the global maxima. The smoothness condition, in particular, is defined as Lipschitz continuity with respect to some metrics (Kleinberg (2005); Kleinberg et al. (2008)) or dissimilarity functions (Kleinberg et al. (2008); Rubbeck et al. (2010)), or $\alpha$-Hölder continuity with $0 < \alpha \leq 1$ (Kleinberg (2005); Auer et al. (2007)). Worst-case lower bound under the Lipschitz assumption is presented in Kleinberg et al. (2008) and that under the Hölder continuity assumption in Locatelli and Carpentier (2018).

Existing works rarely consider the generalization to Hölder space. Recently Hu et al. (2020) studied contextual bandit with reward functions in Hölder spaces, however, the reward function is assumed to be smooth with respect to the observed contexts and the action set is finite. For non-contextual continuum-armed bandits, Akhavan et al. (2020) focus on the strongly convex subset of functions in Hölder spaces with $\alpha \geq 2$ by using projected gradient-like algorithms. Grant and Leslie (2020) analyze Thompson sampling (TS), a Bayesian method, on Hölder spaces with integer-valued exponents and derive a suboptimal upper bound based on the complexity of the function space.

Adaptivity to Smoothness of the Reward Functions. An intriguing problem is whether an algorithm that is oblivious to the Hölder exponent $\alpha$ can simultaneously achieve minimax rates for a range of values for $\alpha$. For non-contextual continuum-armed bandits, this has been discussed only under the Hölder continuous($\alpha \leq 1$) setting. In particular, Locatelli and Carpentier (2018) state that generally, such minimax adaptation to $\alpha$ is impossible by providing a worst-case lower bound for adaptation between two Hölder-continuous function spaces. Additionally, they propose conditions under which it would become possible. (For the contextual finite-armed bandit studied in Hu et al. (2020), Gur et al. (2019) provide lower bounds with similar rates and the extra conditions as well.) However, it remains unclear that, without the extra conditions, whether an algorithm can achieve the lower bound when adapting to a continuous scale of general Hölder spaces.

1They comment that the reason could be either the analysis being suboptimal or the nature of TS. They also derive lower bounds under one-dimension setting, but as we later remark in this paper, the same lower bound can be implied by Wang et al. (2018) under a more general setting.
Another relevant line of work is more broadly model selection in bandit settings, which we will leverage in bandit optimization of Hölder-smooth functions as well as adaptation to the smoothness. In this problem, given a set of base algorithms on possibly different domains, the learner needs to adapt to the best one in an online fashion. The goal is to achieve cumulative regret comparable to the best base algorithm if it were run solely. Bubeck et al. (2011) study the model selection problem for adapting to the unknown Lipschitz constant of functions. Foster et al. (2019) study adapting to the unknown policy dimension in contextual linear bandits by estimating the gap between two policy classes. Agarwal et al. (2016) develop a general algorithm named Corral for bandit model selection under adversarial feedback. It uses online mirror descent to balance between base algorithms. For stochastic feedback particularly, Pacchiano et al. (2020) modify the Corral algorithm to relax requirements on base-algorithms and improve the result on some problem instances (including the one in Foster et al. (2019)). Another relevant issue addressed in Krishnamurthy et al. (2019) which study contextual continuum-armed bandits with Lipschitz continuous reward functions, is their use of the original Corral algorithm applied with EXP4 for adaptation to unknown Lipschitz constant. UCB-type algorithm for corralling base-algorithms is used in Arora et al. (2020) under the assumption that the base-algorithms are finite-armed, and only one of them has access to the best arm.

1.2 Our Contributions

We study bandit optimization of functions in general Hölder spaces. This paper furthers the previous works in the following two main aspects:

1. We propose a novel class of two-layer bandit algorithms, where a carefully-chosen Meta-algorithm deploys misspecified bandit algorithms as arms. Our algorithms show explicitly how to exploit higher-order smoothness in achieving optimal exploration-exploitation tradeoff. We derive worst-case regret bound for this algorithm that matches the existing lower bound except for log factors, for functions in Hölder space including when \( \alpha > 1 \). Our results bridges the gap between Lipschitz smooth bandits where the Hölder exponent is \( \alpha = 1 \) and infinitely-differentiable problems such as linear bandits where the Hölder exponent is \( \alpha = \infty \).

2. We study adaptation to a sequence of Hölder spaces indexed by a continuous but unknown variable of exponent \( \alpha \). We propose a strategy with theoretical guarantee, which uses the bandit model selection algorithm Corral from Pacchiano et al. (2020) applied with versions of our proposed two-layer algorithms. The derived regret bound is to our knowledge the first result on upper bounds when adapting to a continuous scale of Hölder spaces in continuum-armed bandit optimization.

The rest of this paper is organized as follows: In section 2 we introduce the problem formulation and assumptions. We present the two-layer Meta-algorithms and theoretical guarantees in section 3. In section 4 we study the adaptation to unknown smoothness and conclude the paper in section 5 with some open questions.

2 Problem Formulation

In this paper, we consider bandit optimization of smooth functions in Hölder space \( \sum(\alpha, L) \) with \( \alpha > 1 \). The Hölder space is defined formally in definition 1. Some works also study benign problem instances with additional “growth” conditions than the smoothness to characterize the difficulty of finding global maxima, for improvements in regret bounds. For example, Auer et al. (2007) use a parameter to model the growth rate of Lebesgue measure of the near-optimal arms set as a function of the threshold. The near-optimality dimension in Bubeck et al. (2010) uses packing number but has similar meaning. In this paper we will focus solely on worst-case regret to preserve simplicity and leave adaptation to benign cases as a future direction. The performance of the learner is measured by cumulative pseudo-regret as stated below where \( x^* \in \arg\max_{x \in X} \ f(x) \). Throughout this paper we
will simply refer to the pseudo-regret as regret.

\[ R(T) = \sum_{t=1}^{T} [f(x^*) - f(x_t)]. \]  

(1)

To formally define Hölder spaces, we first introduce some notations. Define the following notions for a vector \( s = (s_1 \ldots s_d) \): let \( |s| = s_1 + \cdots + s_d \), \( s! = s_1! \cdots s_d! \) and \( x^s = x_1^{s_1} \cdots x_d^{s_d} \). And define \( D^s = \frac{\partial^{s_1} \cdots \partial^{s_d}}{\partial x_1^{s_1} \cdots \partial x_d^{s_d}} \).

**Definition 1** (Tsybakov (2008)). The Hölder space \( \sum(\alpha, L) \) on domain \( \mathcal{X} \subseteq \mathbb{R}^d \) is defined as the set of functions \( f : \mathcal{X} \to \mathbb{R} \) that are \( l = \lfloor \alpha \rfloor \) times differentiable and have continuous derivatives. \( l \) is the largest integer that is strictly smaller than \( \alpha \). A function \( f \) in \( \sum(\alpha, L) \) satisfies the following inequality for all \( x,y \in \mathcal{X} \).

\[ D^s f(x) - D^s f(y) \leq L \|x - y\|_\infty^{\alpha - l}, \quad \forall s \text{ s.t. } |s| = l. \]

In particular, a function in \( \sum(\alpha, L) \) is close to its Taylor approximation:

\[ |f(x) - T^d_y(x)| \leq L \|x - y\|_\infty^\alpha, \forall x, y \in \mathcal{X}. \]

We use \( T^d_y \) to denote the \( l \)-degree Taylor polynomial around \( y \), \( T^d_y(x) = \sum_{|s| \leq l} \frac{(x-y)^s}{s!} D^s f(y). \)

**Assumptions** We specify the assumptions that are used throughout this paper.

**G1.** The input domain \( \mathcal{X} \) is a hypercube \([0,1]^d\). For simplicity assume the reward function is bounded: \( \|f\|_\infty \leq 1 \).

**G2.** The function \( f \) belongs to Hölder space \( \sum(\alpha, L) \) with some constant \( L > 0 \).

**G3.** The observations are noisy: \( y = f(x) + \eta \) where the noise \( \eta \) is drawn from i.i.d zero mean sub-gaussian distribution with parameter \( \sigma \).

### 3 Meta-algorithm and Analysis

A commonly used method for continuum-armed bandits is fixed discretization, which divides the continuous input domain into finite number of bins, to transform the problem into finite-armed bandit. Previous works mostly consider Hölder-continuous (\( \alpha \leq 1 \)) functions. For example, Auer et al. (2007) study the \( \alpha \)-Hölder continuous functions with \( \alpha \leq 1 \) for one-dimension domain, followed by Bubeck et al. (2010) who generalize it to \( d \)-dimensional domain and propose the HOO algorithm with adaptive discretization. In these works, it suffices to perform random sampling (Auer et al., 2007; Bubeck et al., 2010) or midpoint sampling (Kleinberg, 2005) inside each bin. The worst-case regret bound for Lipschitz space of \( \tilde{O}(T^{\frac{d+\alpha}{d+2}}) \) are matched by the general lower bound of \( \Omega(T^{\frac{d+\alpha}{d+\alpha+1}}) \) (Auer et al., 2007; Bubeck et al., 2010; Locatelli and Carpentier, 2018; Bubeck et al., 2011) apart from log factors. However, if we apply the same methods of random sampling on fixed discretization (Auer et al., 2007) on functions with Hölder exponent \( \alpha > 1 \), the regret incurred is \( \tilde{O}(T^{\frac{d+1}{d+2}}) \) since the Hölder space with exponent \( \alpha > 1 \) is a subset of the Lipschitz function space. It prompts us to ask the question of whether a better rate that matches the dependence on \( \alpha \) can be achieved for functions that are smoother than Lipschitz. An extreme is when \( \alpha \) reaches infinity, where the reward model will be infinitely-differentiable, for example the stochastic linear bandit which enjoys \( \tilde{O}(T^{\frac{1}{2}}) \) regret even on continuous domain (Dani et al., 2008; Abbasi-Yadkori et al., 2011).

\(^2\)Only when referring to the order of Hölder smooth functions’ derivatives do we denote \( \lfloor \cdot \rfloor \) as the largest integer strictly less than input. In other places in this paper it denotes less or equal to input.

\(^3\)We use \( l_\infty \) norm as in some works on adaptive confidence bands and optimization (Low et al., 1997; Tsybakov, 2008; Hofmann et al., 2011; Wang et al., 2018).

\(^4\)In this paper, for simplicity, we assume \( L \) is some constant that satisfies assumption G1.

\(^5\)The adaptive discretization does not change worst-case regret but has improvements on benign problems, as introduced in section 2.
3.1 Algorithm Overview

We keep to fixed discretization of the domain since we consider only the worst-case regret. We divide \( \mathcal{X} = [0, 1]^d \) into \( n \) equal-sized hypercubes, leaving \( n \) as a parameter of the algorithm. As shown in definition\(^1\) the function is locally well-approximated by Taylor polynomial which reduces to a linear model of a feature map of \( x \) with dimension \( d(\alpha) \). It is equivalent to observing a misspecified linear model inside each bin, the equivalence formally quantified in Lemma\(^2\). Therefore, local exploration-exploitation tradeoff can be achieved by a base algorithm with sublinear regret on such misspecified models, with a Meta-algorithm to balance the budgets between the base algorithms in the bins.

**Lemma 2.** Let hypercube \( B_\Delta \) be a subset of the input space with volume \( \Delta \). If a function satisfies assumption G1 \( \sim 2 \), there exists a linear parameter \( \theta^* \in \mathbb{R}^{d(\alpha)} \) and feature map \( \phi : x \mapsto \phi(x) \in \mathbb{R}^{d(\alpha)} \), such that \( f \) can be approximated by the linear function: \( \| f - (\theta^*, \phi(x)) \|_\infty \leq \epsilon = L\Delta^{\frac{2}{3}} \) for \( x \in B_\Delta \). When \( \alpha \leq 2 \), \( d(\alpha) = d \); when \( \alpha > 2 \), \( d(\alpha) = \mathcal{O}(d^l) \) with \( l \) (definition\(^1\)). Note that the linear parameter may not be unique.

The proof is in Appendix section\[^{A.1} \]. In the following parts of this section we first present the misspecified bandit algorithm to run inside a bin, and then the Meta-algorithms to control these local algorithms.

3.2 The Misspecified Linear Bandit Algorithm

In this subsection we escape from the big picture briefly in order to present the misspecified linear bandit algorithm, modified from the ConfidenceBall\(_2\) algorithm in Dani et al. (2008) to serve as “arms” of the Meta algorithm. The algorithm, as shown in its name, is based on construction of confidence ellipsoid of the unobserved linear parameter in dimension \( d \). We prove that the proposed modification can accommodate bias in the function feedback by deriving an upper bound on the cumulative regret of \( \mathcal{O}(d\sqrt{T} +dT\epsilon) \). Here \( \epsilon \) is the upper bound on bias value and known by the algorithm. We recently discovered that a similar result with proof sketch already appeared in recent work of Lattimore and Szepesvari (2019) (appendix E) who used modification of the algorithm in Abbasi-Yadkori et al. (2011), and hence enjoys the improvement of a multiplicative factor \( \sqrt{\log(T)} \). For completeness and to provide necessary intermediate results for Meta-algorithms in later sections, we present our algorithm and full proof as complementary. It is worth mentioning that without the modification, the original algorithm incurs suboptimal regret under misspecification.

**Assumptions** We make the following assumptions for the misspecified model. Note that they are consistent with the aforementioned global assumptions.

**A1.** The feedback model is \( y = \langle x, \theta^* \rangle + b(x) + \eta \) with \( |b(x)| \leq \epsilon, \forall x \in \mathcal{X} \in \mathbb{R}^d \).

**A2.** The mean reward \( \mathbb{E}[y] \) is bounded by \([-1, 1]\).

**A3.** The noise \( \eta \) is drawn from zero-mean sub-gaussian with parameter \( \sigma \).

The pseudo-code of the modified algorithm is shown in Algorithm\[^1\]. The goal is to minimize the cumulative pseudo-regret of the linear model:

\[
R(T) = \sum_{t=1}^{T} r_t = \sum_{t=1}^{T} (\langle x^*, \theta^* \rangle - \langle x_t, \theta^* \rangle).
\]

We prove that this regret is \( \mathcal{O} \left( d\ln(T)\sqrt{\ln\left(\frac{T^2}{\delta}\right)}T + cTd\sqrt{2\ln(T)} \right) \) with probability \( 1 - \delta \). This is formally stated in Theorem\[^3\].

\[^{a}\]We slightly abuse the notation and define short-hand notation \( \langle \theta, x \rangle := \theta_0 + \sum_{i=1}^d \theta_i x_i \).

\[^{b}\]For clarity this use of \( \mathcal{O} \) omits \( \ln(T) \) and \( \delta \) dependence.

\[^{c}\]Different from Dani et al. (2008) who assumes bounded noise. This reflects in the difference in \( \beta_t \).
We now present the first structure of our Meta-algorithms. We consider the most straightforward Algorithm 1

Algorithm 1: Misspecified linear UCB algorithm ($A^{\text{local}}$)

**Require:** Misspecification error $\epsilon$, input domain $\mathcal{X}$ and its dimension $d$, fail probability $\delta$, upper bound on $\|x\|_2^2$, $\kappa^2 = d$.

1: Initialize $A_1 = I_d$ and $x_1 \in \mathcal{X}$ randomly.
2: for $t = 1 \ldots$ do
3: Execute action $x_t$ and observe reward $y_t$
4: $A_{t+1} = A_t + x_tx_t^T$
5: $\theta_{t+1} = A_{t+1}^{-1}(\sum_{r=1}^{t} y_rx_r)$
6: $\beta_{t+1} = 128\sigma^2d\ln(1+t)\ln(4\frac{(t+1)^2}{\delta})$
7: Define function $UCB_{t+1}(x) = \left(\langle x, \hat{\theta}_{t+1} \rangle + \sqrt{\beta_{t+1} \|A_{t+1}^{-1/2}x\|} + \epsilon \sum_{s=1}^{t} |x^T A_{t+1}^{-1} x_s| \right)$
8: Compute action $x_{t+1} = \arg\max_{x \in \mathcal{X}} UCB_{t+1}(x)$
9: Return $x_{t+1}$ and $UCB_{t+1}(x_{t+1})$
10: end for

**Theorem 3.** If assumptions A1~A3 hold, then with probability $1 - \delta$, the cumulative regret of Algorithm 1 is upper bounded by:

$$R(T) \leq \sqrt{8d\beta_T T \ln(1+T)} + 2\epsilon T d \sqrt{2 \ln(1+T)} + 2\epsilon T.$$  

The first term is the standard stochastic linear bandit regret rate same as in Dani et al. (2008). We defer the proof to Appendix section A.2. The increment of a multiplicative factor $\sqrt{d}$ in the second term compared to that in Lattimore and Szepesvari (2019) is due to difference in assumption on $\|x\|^2$. Their assumption is $\|x\|^2 \leq 1$ whereas ours is $\|x\|^2 \leq d$.

### 3.3 The UCB-Meta-algorithm

We now present the first structure of our Meta-algorithms. We consider the most straightforward structure: UCB-Meta, the pseudo-code is shown in Algorithm 2 (define $\lfloor \cdot \rfloor$ as the action of rounding to nearest integer). We keep a version of the base mispecified linear bandit algorithm in each bin.

The confidence estimates of the local linear models are passed to the Meta-algorithm as UCB of their assumption is

$$\|x\|^2 \leq 1$$

of the confidence bands as well as their lengths. In particular, if the function value $f(x)$ at time $t$ is contained in an honest confidence band $[UCB_t(x) - 2l_t(x), UCB_t(x)]$, then we can use the length $l_t(x)$ to bound instantaneous regret incurred by the selected action at this step. The confidence ellipsoids for the piecewise linear parameters $\hat{\theta}_{k,t}$ that are constructed by local mispecified linear bandits offer a convenient confidence estimation of function value, with the additional adjustment factor $\epsilon$, the approximation error. The full proof is deferred to Appendix section A.3. The algorithm

```
for $t = 1 \ldots$ do
  Execute action $x_t$ and observe reward $y_t$
  $A_{t+1} = A_t + x_tx_t^T$
  $\theta_{t+1} = A_{t+1}^{-1}(\sum_{r=1}^{t} y_rx_r)$
  $\beta_{t+1} = 128\sigma^2d\ln(1+t)\ln(4\frac{(t+1)^2}{\delta})$
  Define function $UCB_{t+1}(x) = \left(\langle x, \hat{\theta}_{t+1} \rangle + \sqrt{\beta_{t+1} \|A_{t+1}^{-1/2}x\|} + \epsilon \sum_{s=1}^{t} |x^T A_{t+1}^{-1} x_s| \right)$
  Compute action $x_{t+1} = \arg\max_{x \in \mathcal{X}} UCB_{t+1}(x)$
  Return $x_{t+1}$ and $UCB_{t+1}(x_{t+1})$
end for
```
Algorithm 2 UCB-Meta-algorithm ($A_{\text{global}}$)

Require: smoothness parameter $\alpha$, Hölder constant $L$, dimension of domain $d$, time horizon $T$ and fail probability $\delta$, action space $X$.

1: Initialize $n = \lceil T^{\frac{d+\alpha}{2d}} / \ln(T)^{\frac{d+\alpha}{2d}} \rceil$ and divide the action space $X$ into same-sized bins $B_1...n$ with volume $\Delta = 1/n$.
2: for $k = 1, \ldots, n$ do
3: On bin $B_k$, start a version of local misspecified base-algorithm $A_k$ using misspecification error $\epsilon = \ln \frac{2}{\delta}$, input domain $X^* = \{\phi(x), x \in X\}$ and its dimension $d(\alpha)$, fail probability $\delta/n$.
4: Initialize counter $s_k = 1$ to indicate how many times $A_k$ is queried.
5: Query $A_k$ once by running steps 3-9 of Algorithm 1 with $t = s_k$ and obtain upper confidence bound $\text{UCB}_k$.
6: $s_k \leftarrow s_k + 1$
7: end for
8: for $\tau = 1 \ldots T$ do
9: Select the bin with index $k(\tau) = \arg\max_k \text{UCB}_k$.
10: Execute the local bandit algorithm $A_{k(\tau)}$ once by running steps 3-9 (of Algorithm 1) with $t = s_{k(\tau)}$.
11: Receive updated recommendation $\phi_\tau \in \{\phi(x), x \in B_{k(\tau)}\}$ and $\text{UCB}_{k(\tau)}$.
12: Advance counter for $A_{k(\tau)}$: $s_{k(\tau)} \leftarrow s_{k(\tau)} + 1$.
13: end for

defines each bin to be a hypercube with volume $\Delta = 1/n$, according to Lemma 2 we have $\epsilon = \ln \frac{2}{\delta}$. Therefore, setting $n = \mathcal{O}(T^{\frac{d+\alpha}{2d}} / \ln(T)^{\frac{d+\alpha}{2d}})$ will minimize the upper bound and yield cumulative regret bound of $R(T) \leq \tilde{\mathcal{O}}(d(\alpha)T^{\frac{d+\alpha}{2d}})$. (5)

3.3.2 Anytime Regret Guarantee for Algorithm 2

To achieve the rate in bound 5, Algorithm 2 needs to know the time horizon $T$ in advance to set $n$ and $\epsilon$ correspondingly. Here we prove that, with the doubling trick (Auer et al. (1995)), the UCB-Meta-algorithm can get regret that is of the same rate as in bound 5 up to constant factors without knowing $T$. This result is needed in the adaptation problem studied in section 4.

Theorem 5. If Algorithm 2 with access to the time horizon $T$ achieves regret of $\tilde{\mathcal{O}}(T^a)$ with probability $1 - \delta$, then the procedure described in Algorithm 3 can achieve regret rate $\tilde{\mathcal{O}}(T^a)$ with probability $1 - \delta$ without the knowledge of $T$.

The pseudo-code for Algorithm 3 is in Appendix section B.1 and the proof of Theorem 5 in Appendix A.3.4.

3.4 The Corral-Meta-algorithm

Another choice for Meta-algorithm is bandit model selection methods. Here we use the Corral algorithm defined in [Pacchiano et al. (2020)], which will be introduced more formally in section 4. An example of corraling misspecified linear bandit algorithms without corruption to the regret rate apart from log factors has already been given in [Pacchiano et al. (2020)], but for adaptation to the misspecification error $\epsilon$. Here we demonstrate that it can also be used to corral misspecified bandit base-algorithms on different bins in a discretized domain. We derive the following regret bound that is the same as UCB-Meta-algorithm.

Theorem 6. First perform the smoothing transformation (Algorithm 3 in [Pacchiano et al. (2020)]) to our misspecified linear bandits in Algorithm 7, denote the smoothed misspecified linear bandits
as $A^\text{local}_i$. Then, the Meta-algorithm (Algorithm 5 (Corral-Update) reproduced in Pacchiano et al. (2020)) applied with a set of $A^\text{local}_i$ that are initialized in the same way as in Algorithm 2 has expected regret upper bounded by:
\[
\mathbb{E}[R(T)] \leq \tilde{O}(d(\alpha)T^{\frac{d+\alpha}{\alpha+2}}).
\] (6)

The proof of this theorem is in Appendix section A.4.

### 3.5 Discussion

The role of the Meta-algorithm is essentially model selection and adaptation to the base-algorithms. It is not a trivial task since the rewards incurred by the base-algorithms are not i.i.d as in standard stochastic settings. However, UCB as a stochastic multi-armed bandit algorithm, is applicable as Meta-algorithm because the local parametric (linear) function approximations provide honest upper confidence bounds for each bin even under the misspecifications, thus enabling the distribution-independent analysis for UCB. The advantage of Corral-Meta is that it potentially allows relaxation of the Hölder smoothness to hold only around the global maxima (Auer et al. (2007); Bubeck et al. (2010)), while the same relaxation is not straightforward for UCB-Meta. The advantage of UCB is that under standard stochastic settings where each arm has i.i.d rewards, it achieves the gap-dependent bound of $O(\log(T)/\Delta)$. Thus an interesting question for the future is whether similar gap-dependent bounds for the UCB-Meta is available. Such bounds would enable exploitation of the growth conditions (section 2) for potential rate improvements.

### 3.6 Comparison with Existing Lower Bound

We compare the derived upper bounds of $\tilde{O}(d(\alpha)T^{\frac{d+\alpha}{\alpha+2}})$ to the existing lower bound from Wang et al. (2018), which study global optimization. In their work, the performance of optimization algorithms with output $\hat{x}_T$ is measured by simple regret $\mathcal{L}(\hat{x}_T; f) \triangleq f(\hat{x}_T) - f(x^*)$, for $f$ in Hölder spaces including $\alpha \geq 1$. Theorem 2 (coupled with Proposition 3) in Wang et al. (2018) implies that $\sup_{f \in \sum(\alpha)} \mathbb{E}[\mathcal{L}(\hat{x}_T; f)] = \Omega(T^{\frac{\alpha}{\alpha+2}})$. We argue that this lower bound can be directly used to lower bound the worst-case cumulative regret, by making the following observation (remark 3 in Bubeck et al. (2010)): If a strategy achieves expected cumulative regret $\mathbb{E}[R_T]$, then by uniformly selecting a past action as the final output $\hat{x}_T$, it can also achieve expected simple regret $\mathbb{E}[\mathcal{L}(\hat{x}_T; f)] = \mathbb{E}[R_T]/T$. Therefore, any strategy with cumulative regret $\tilde{o}(T\mathbb{E}[\mathcal{L}(\hat{x}_T; f)])$ will violate the lower bound. Through proof by contradiction, we take the result from Wang et al. (2018) as an $\Omega(T^{\frac{\alpha}{\alpha+2}})$ lower bound on expected cumulative regret, and argue that our results match this bound up to log factors. Our results show that proposed algorithms are minimax optimal in dependence of $T$ and effectively exploit the function smoothness.

### 4 Adaptation to Unknown Smoothness

In this section, we study adaptation to the smoothness exponent $\alpha$ of the reward function. Minimax adaptation, which means a learner can simultaneously achieve the minimax optimal rates (Hoffmann et al. (2011); Locatelli and Carpentier (2018)) under a nested set of Hölder spaces, has been proven to be impossible for cumulative regret minimization without additional assumptions. Locatelli and Carpentier (2018) provide a lower bound for adaptation between two Hölder continuous functions spaces. Assume $\alpha < \gamma \leq 1$, for any strategy with a good expected regret $\mathbb{E}[R_\gamma(T)]$ in $\sum(\gamma, L)$, they show that its expected regret in the superset $\sum(\alpha, L)$ will depend inversely on $\mathbb{E}[R_\gamma(T)]$, and therefore be suboptimal for $\sum(\alpha, L)$. They propose a strategy to match that lower bound that requires values of $\alpha$ and $\gamma$, thereby also proving that the lower bound is tight.

However, when adapting to a continuous scale of Hölder spaces (possibly $\alpha \geq 1$), it remains unclear what strategy can generalize and achieve this lower bound for some Hölder spaces. We aim to answer that question by proposing a new strategy that uses a recently developed bandit model selection algorithm (Corral with smooth wrapper in Pacchiano et al. (2020)) applied with a set
of Meta-algorithms (section 3). We will present this strategy and its theoretical guarantees next. Throughout the following sections, we refer to minimax optimal in dependence of T as minimax unless otherwise specified.

## 4.1 Corral Applied with Meta-algorithms

The bandit model selection method Corral is first developed by Agarwal et al. (2016) and based on an instance of online mirror descent with mirror map derived from Foster et al. (2016). Corral with smooth wrapper proposed by Pacchiano et al. (2020) for stochastic feedback problems is different from the original Corral algorithm in the following aspects. The smoothed version no longer needs to send importance-weighted feedback to base-algorithm, therefore no longer requires the base-algorithms themselves to be modified for stability guarantee (definition 3 in Agarwal et al. (2016)). In the following parts, we will use Corral with smooth wrapper to adapt to the smoothness and refer to it as Corral for simplicity. A copy of the pseudo-code of Corral from Pacchiano et al. (2020) can be found in Appendix B.2 for easier reference. We use a set of $M$ Meta-algorithms $A^\alpha$ as bases. The input values $\alpha_i$ are from a grid $G$ defined later. Therefore, we first specify the regret of a Meta-algorithm with input smoothness parameter $\alpha'$ that is ran on functions with actual Hölder smoothness $\alpha$.

**Lemma 7.** For function $f$ that satisfies global assumptions $G1 \sim G3$ with parameter $\alpha$, the regret of Algorithm 2 with input parameter $\alpha' \leq \alpha$ is bounded with probability $1 - \delta$ by

$$R(T) \leq \hat{O}(d(\alpha')T^{\frac{d + \alpha'}{2 + \alpha'}}).$$

(7)

The bound does not hold for $\alpha' > \alpha$.

The proof is deferred to Appendix section A.3. Having established the performance of base algorithms with misspecified smoothness exponents, we present the adaptation strategy and its regret bound in Theorem 8. Since it is impossible to achieve minimax optimal rates for multiple values of the smoothness parameter simultaneously, we introduce a user-sepecified parameter $R$ that controls the Hölder space over which minimax optimality is desired. We show that conditioned on achieving minimax rate for the space $\sum (R, L)$, our adaptation strategy provides best possible regret bound on all supersets $\sum (\alpha, L)$ where $\alpha \leq R$. The results are stated in Theorem 8.

**Theorem 8.** Consider adapting to a continuous scale of nested Hölder spaces indexed by $\alpha$ whose value is bounded in a given interval, for simplicity we assume $0 < \alpha < 2$, where $d(\alpha) = d$. Define $R \leq 2$ as a parameter set by the decision-maker that specifies the index of Hölder space for which minimax optimal regret is achieved. Define linear grid $G = \{\alpha_i = \frac{R}{\log(T)}i, i = 0, 1 \ldots \lfloor \log(T) \rfloor\}$ so that the total number of base algorithms is $M = |G| = \lfloor \log(T) \rfloor$. Consider using Corral with bases that are Meta-algorithms (algorithm 2 in Appendix section B.1) with input $\alpha_i \in G, i \in [M]$. Then by setting the learning rate of Corral to be $\eta = d^{-1}T^{-\frac{d}{d + R}}$, the regret rates achieved for any Hölder exponent $\alpha \in (0, 2]$ are:

$$\sup_{f \in \sum(\alpha, L)} \mathbb{E}[R(T)] \leq \hat{O}(dT^{\frac{d^2 + 2Rd + R^2}{d + R + 2d(\alpha d + \alpha)}}) \text{ for } \alpha \in (0, R],$$

(8)

$$\sup_{f \in \sum(\alpha, L)} \mathbb{E}[R(T)] \leq \hat{O}(dT^{\frac{d + R}{d + R + 2d(\alpha d + \alpha)}}) \text{ for } \alpha \in [R, 2].$$

(9)

A straightforward example is shown in Figure 11. Functions with Hölder exponent $\alpha > R$ essentially belongs to a subset of $\sum (R, L)$ and have the same regret rates as in equation (9) because the algorithm did not fully exploit their smoothness. There are two sources of cost of adaptation, first the cost of adapting to $M$ grid points. Since $M = O(\log(T))$, this has the same difficulty as the adaptation to two values in Locatelli and Carpentier (2018). The second one, however, is a consequence of adapting to a continuous scale of $\alpha$. The cost is the rate difference between the exponent $\alpha$ and the closest value to it on $G$, denoted $\hat{\alpha} \in G$, s.t. $\hat{\alpha} \leq \alpha \leq \hat{\alpha} + \frac{R}{\log(T)}$. This cost can be alleviated by the design of the linear grid. We defer the full proof to Appendix section A.6.

---

11Since the core of oneline mirror descent in Corral is not changed.
4.2 Comparison with Existing Lower Bound for Adaptation

In this subsection, we compare the results in Theorem 8 to the existing lower bound in Locatelli and Carpentier (2018). Theorem 3 of Locatelli and Carpentier (2018) states that given two smoothness values $\alpha_1 < \alpha_2 \leq 1$, if a strategy has expected regret $\mathbb{E}[R_{\alpha_2}(T)]$ under exponent $\alpha_2$ that is $\tilde{O}(T d + R d^{\alpha_2/2})$, then the regret of this strategy under the superset characterized by $\alpha_1$ is lower bounded by

$$\sup_{f \in \Sigma(\alpha_1, L)} \mathbb{E}[R(T)] \geq \tilde{\Omega}(T R_{\alpha_2}(T)^{-\alpha_1/\alpha_2})$$

even if the strategy has access to both $\alpha_1$ and $\alpha_2$.

We make the following remark: for any pair of exponent values $(\alpha_1, \alpha_2)$ where $\alpha_1 < R$ and $R \leq \alpha_2 \leq 1$, the strategy proposed in Theorem 8 matches the lower bound except for log factors. We verify this by plugging in $\mathbb{E}[R_{\alpha_2}(T)] = \tilde{O}(T d + R d^{\alpha_2/2})$, omitting dependence on $d$, to yield the lower bound on $\Sigma(\alpha_1, L)$ which is $\tilde{O}(T d^{2 \alpha_2 / (\alpha_2 + 2)} + R_{\alpha_2}^{\alpha_1 / \alpha_2})$. This is matched by our upper bound in equation (8), apart from log factors and $d$. An illustration is shown in Figure 2. In other words, the proposed algorithm can perform under unknown smoothness exponent and match the lower bound (available only for exponent values within $(0, 1]$) on a subset of Hölder spaces.

5 Conclusion

The core of this paper is extending the assumption on function space from Lipschitz to Hölder spaces with higher-order smoothness in bandit optimization of black-box functions. We also study adaptation to the smoothness under this scope. The class of two-layer algorithms that we proposed consists of a Meta-algorithm with the choice of UCB (Auer et al. (2002)) or Corral (Agarwal et al. (2016); Pacchiano et al. (2020)) and a set of misspecified bandit base-algorithms as arms. We derive regret upper bounds for $\alpha$-Hölder smooth functions with $\alpha > 1$ that matches existing lower bounds in their dependence on $T$, the number of active queries, with straightforward generalization to larger $\alpha$. Our framework provides useful insights in exploiting higher-order smoothness of reward functions for cumulative regret minimization, because our two-layer structure allows base-algorithms to perform local exploration-exploitation tradeoff as opposed to the local pure exploration done for bandit optimization of $\alpha$-Hölder continuous functions. For adaptation to the smoothness exponent, we further previous works by deriving regret upper bound for adaptation to a continuous scale of Hölder spaces with exponent $\alpha$ in a given range. We show that by using bandit model selection algorithms, it can achieve the existing lower bound between two Hölder spaces, even if the algorithm does not know both exponent values.

Our work inspires several directions for the future. An intriguing direction is to study whether there exist gap-dependent bounds for the UCB-Meta algorithm, whose arms have non i.i.d rewards because they are bandit algorithms themselves. Such bounds could enable better rates for benign problem instances, for example with the growth conditions (mentioned in section 2). Another direction is the
relaxation of the Hölder smooth assumption, to hold only around the maxima instead of everywhere on $\mathcal{X}$, which is considered by prior works such as [Auer et al. (2007); Kleinberg et al. (2008); Bubeck et al. (2010)]. Finally, it remains an open problem to establish the lower bound for adaptation when the smoothness exponents are larger than 1.

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Appendix of Smooth Bandit Optimization: Generalization to Hölder Space

A Auxiliary proofs for the main document

A.1 Proof of Lemma 2

Proof. Recall the definition of Hölder smoothness: \( |f(x) - T^l_y(x)| \leq L\|x - y\|_\infty^\alpha \). For a hypercube \( B, \|x - y\|_\infty \leq \Delta \), \( x, y \in B \). By definition, when the function smoothness exponent \( \alpha \in (1, 2] \), \( l = 1 \). Notice that the Taylor polynomial of degree \( l = 1 \) around \( y \) is a linear function of \( x \): 
\[
T^1_y(x) = f(y) + \frac{\partial f}{\partial x}(y)(x - y) + \frac{\partial^2 f}{\partial x^2}(y)(x - y)^2 + \ldots
\]
When \( \alpha \geq 2 \), the Taylor polynomial still be written as a linear function but of higher-dimensional feature map of \( x \): \( \phi : [0, 1]^d \to [0, 1]^{d(\alpha)} \) which contains exponentiations of elements in \( x \). Using the operations defined for definition \( \phi(x) = \{x^s, \forall s, s.t. |s| \leq l\} \) so:
\[
d(\alpha) = |\{s : 1 \leq |s| \leq l\}| = \sum_{1 \leq j \leq l} \left( j + d - 1 \right) d - 1 = O(d^l)
\]
When \( l = 1 \), it is equivalent to defining \( \phi(x) = x \). The parameter \( \theta \) is determined by the derivatives of \( f \) at \( y \) and the value of \( y \). Therefore, we know locally there exists a linear parameter in dimension \( \theta^* = \arg \min_y \|f - \phi(x)^T \theta\|_\infty, x \in B \), such that \( \|f - \langle \theta^*, \phi(x) \rangle\|_\infty \leq \epsilon = L\Delta \), \( \forall x \in B \). Also, note that \( \|\phi(x)\|_2^2 \leq d(\alpha)^2 \) according to definition. When the exponent \( \alpha \in (0, 1] \), \( l = 0 \) and the Taylor polynomial is simply a constant. Therefore the same argument holds for \( \theta^* \) for example when \( \theta^1, \ldots, \theta^d = 0 \) (a constant function).

A.2 Proof of Theorem 3

Proof. Throughout this proof, we assume that the assumptions A1~3 hold. This proof is modified from that in [Dani et al., 2008]. Some techniques are from [Abbasi-Yadkori et al., 2011]. We only present the parts which we change. First we proof the following bound on simple regret at each step:
\[
r_t \leq 2\sqrt{\beta_t}||A_t^{-1/2}x|| + 2\epsilon \sum_{s=1}^{t-1}||x^T A_t^{-1} x_s||.
\]
(11)

And then we will bound the sum of these two terms separately. In order to proof inequality \( |\{s : 1 \leq |s| \leq l\}| = O(d^l) \) we start from an important auxiliary theorem of confidence bound on \( \theta^* \), Theorem 9.

Theorem 9. Let \( \beta_t = C\sigma^2 d \ln(1 + t\kappa^2/d) \ln(2d^2) \left( = O(d \ln(t) \ln(d^2)) \right) \) for a sufficiently large constant \( C \), then with probability \( 1 - \delta \), \( \theta^* \) is contained in the confidence set:
\[
\tilde{C}_t = \{\hat{\theta}_t + \sqrt{\beta_t} A_t^{-1/2} z_d - A_t^{-1} (\sum_{s=1}^{t-1} b_s x_s), \|z_d\|_2 \leq 1\},
\]
and as a result,
\[
|\langle x, \theta^* \rangle| \leq |\langle x, \hat{\theta}_t \rangle| + \sqrt{\beta_t} ||A_t^{-1/2} x|| + \epsilon \sum_{s=1}^{t-1} |x^T A_t^{-1} x_s|,
\]

We slightly abuse the notation and define short-hand notation \( \langle \theta, x \rangle := \theta_0 + \sum_{i=1}^{d(\alpha)} \theta_i x_i \).
The proof of Theorem 9 is in Appendix A.2.1. Now, if \( \theta^* \in \tilde{C}_t \), we have

\[
\begin{align*}
    r_t &= \langle x^*, \theta^* \rangle - \langle x_t, \theta^* \rangle \\
    &\leq \langle x^*, \theta^* \rangle - UCB_t(x^*) + UCB_t(x_t) - \langle x_t, \theta^* \rangle \\
    &\leq UCB_t(x_t) - \langle x_t, \theta^* \rangle \\
    &\leq 2\sqrt{\beta_t} \|A_t^{-1/2}x_t\| + 2\epsilon \sum_{s=1}^{t-1} |x_t^TA_t^{-1}x_s|.
\end{align*}
\]

The first inequality is because our algorithm will only choose \( x_t \) when \( UCB_t(x_t) \geq UCB_t(x^*) \). The last inequality holds because

\[
\begin{align*}
    \langle x, \theta^* \rangle &\geq \langle x, \hat{\theta}_t \rangle + \min_{z_d \in B_d^2} \sqrt{\beta_t} \langle x, A_t^{-1/2}z_d \rangle - \sum_{s=1}^{t-1} b_s x_s^T A_t^{-1} x_s \\
    &\geq \langle x, \hat{\theta}_t \rangle - \sqrt{\beta_t} \|A_t^{-1/2}x_t\| - \sum_{s=1}^{t-1} b_s x_s^T A_t^{-1} x_s \\
    &\geq UCB_t(x) - 2\sqrt{\beta_t} \|A_t^{-1/2}x_t\| - 2\epsilon \sum_{s=1}^{t-1} |x_s^T A_t^{-1} x_t|.
\end{align*}
\]

By assumption on the mean reward function value, the absolute value of instant pseudo-regret \( |r_t| \) is bounded by \( 1 + \epsilon \). Therefore, combining inequality (11) and \( r_t \leq 2 + 2\epsilon \), we have that

\[
r_t \leq (2 + 2\epsilon) \land \left( 2\sqrt{\beta_t} \|A_t^{-1/2}x_t\| + 2\epsilon \sum_{\tau=1}^{t-1} \|x_{\tau}^T A_{\tau}^{-1} x_{\tau}\| \right) \\
\leq 2 \left( 1 \land \sqrt{\beta_t} \|A_t^{-1/2}x_t\| \right) + 2\epsilon \sum_{\tau=1}^{t-1} \|x_{\tau}^T A_{\tau}^{-1} x_{\tau}\| + 2\epsilon. \tag{12}\]
\]

Sum of term #1 is bounded using bound (28) and Cauchy Schwartz inequality:

\[
2\sum_{t=1}^{T} (1 \land \sqrt{\beta_t} \|A_t^{-1/2}x_t\|) \leq 2 \sqrt{T \beta T} \sum_{t=1}^{T} (1 \land \|x_t^T A_t^{-1} x_t\|) = \sqrt{8d \beta T \ln(1 + T \kappa^2/d)}. \tag{13}\]

For sum of term #2, we first have

\[
\begin{align*}
    \sum_{\tau=1}^{t-1} x_{\tau}^T A_{\tau}^{-1} x_{\tau} &\leq \sum_{\tau=1}^{t-1} x_{\tau}^T A_{\tau}^{-1} x_{\tau} x_{\tau}^T A_{\tau}^{-1} x_{\tau} \\
    &= t x_{t}^T A_{t}^{-1} \left( \sum_{\tau=1}^{t-1} x_{\tau} x_{\tau}^T \right) A_{t}^{-1} x_{t} \\
    &\leq t x_{t}^T A_{t}^{-1} \left( \sum_{\tau=1}^{t-1} x_{\tau} x_{\tau}^T + I_d \right) A_{t}^{-1} x_{t} \\
    &= \sqrt{t x_{t}^T A_{t}^{-1} \left( \sum_{\tau=1}^{t-1} x_{\tau} x_{\tau}^T + I_d \right) A_{t}^{-1} x_{t}}.
\end{align*}
\]

\(^{13}a \land b = \min(a, b)\)
Then the sum \( \sum_{t=1}^{T} (\sum_{\tau=1}^{t-1} x_{t}^{T} A_{t}^{-1} x_{\tau}) \) can be bounded by:

\[
\sum_{t=1}^{T} (\sum_{\tau=1}^{t-1} x_{t}^{T} A_{t}^{-1} x_{\tau}) \leq \sum_{t=1}^{T} (\sqrt{t x_{t}^{T} A_{t}^{-1} x_{t}})
\leq \sqrt{(\sum_{t=1}^{T} t)(\sum_{t=1}^{T} x_{t}^{T} A_{t}^{-1} x_{t})}.
\]

Now, we need to bound \( \sum_{t=1}^{T} x_{t}^{T} A_{t}^{-1} x_{t} \) with inequality (28). We know that \( A_{t}^{-1} \) is a full-rank matrix. Therefore, denote its eigenvalues and eigenvectors as \( \lambda_{1} \ldots \lambda_{d}, v_{1} \ldots v_{d} \). Then

\[
x_{t}^{T} A_{t}^{-1} x_{t} = (c_{1}v_{1} + \cdots + c_{d}v_{d})^{T} A_{t}^{-1} (c_{1}v_{1} + \cdots + c_{d}v_{d})
= c_{1}^{2}\lambda_{1} + \cdots + c_{d}^{2}\lambda_{d}
\leq \lambda_{\max}(A_{t}^{-1})\|x_{t}\|_{2}^{2} = \frac{\kappa^{2}}{\lambda_{\min}(A_{t})}
\leq \frac{\kappa^{2}}{\lambda_{\min}(I_{d}) + \lambda_{\min}(X_{t}^{T} X_{t})} \leq \kappa^{2}.
\]

The second last inequality holds due to Weyl's inequality. Therefore,

\[
\sum_{t=1}^{T} x_{t}^{T} A_{t}^{-1} x_{t} \leq \kappa^{2} \sum_{t=1}^{T} (x_{t}^{T} A_{t}^{-1} x_{t} \wedge 1)
\leq \kappa^{2}(2d\ln(1 + T\kappa^{2}/d)).
\]

Putting the above together,

\[
\sum_{t=1}^{T} \left(2\epsilon \sum_{\tau=1}^{t-1} x_{t}^{T} A_{t}^{-1} x_{\tau}\right) \leq 2\epsilon \sqrt{(\sum_{t=1}^{T} t)(\sum_{t=1}^{T} x_{t}^{T} A_{t}^{-1} x_{t})}
\leq 2\epsilon T\kappa \sqrt{2d\ln(1 + T\kappa^{2}/d)}.
\]

Finally, plugging in \( \kappa^{2} = d \) gives the final results. \( \square \)

### A.2.1 Proof of Theorem 9

**Proof.** Let \( \hat{\theta}_{t} = A_{t}^{-1} X_{t}^{T} y \) denote the regularized least square estimator at time \( t \). Matrix \( X_{t} \) has dimension \((t-1) \times d\), where each row is a past action (until time \( t \)). We first define an unobserved variable \( \tilde{\theta}_{t} \):

\[
\tilde{\theta}_{t} = A_{t}^{-1} X_{t}^{T} (X_{t} \theta^{*} + \eta_{t}) = \hat{\theta}_{t} - A_{t}^{-1} X_{t}^{T} b_{t},
\]

here we abuse the notations and let \( \eta_{t} \) and \( b_{t} \) be the \((t-1) \times 1\) vector containing noise and bias of each time. Then we define the following confidence ellipsoid centered at \( \tilde{\theta}_{t} \):

\[
C_{t} = \{ \theta : (\theta - \tilde{\theta}_{t})^{T} A_{t} (\theta - \tilde{\theta}_{t}) \leq \beta_{t}\},
\]

and prove the following lemma as an analog to Theorem 5 of [Dani et al., 2008]:

**Lemma 10.** The true linear parameter \( \theta^{*} \) is contained in ellipsoid \( C_{t} \), specifically, \( P(\forall t, \theta^{*} \in C_{t}) \geq 1 - \delta \).

\footnote{This proof is extracted from a remark in proof of Theorem 3 in Abbasi-Yadkori et al., 2011}
The proof is in Appendix section A.2.2. However, we do not observe the vector $b_t$, so we cannot calculate $C_t$ in our algorithm. So instead, we define a larger $\tilde{C}_t$ that contains $C_t$, which will naturally contains $\theta^*$ with high probability. To construct $\tilde{C}_t$, we first re-write $C_t$ as

$$C_t = \{ \hat{\theta}_t + \sqrt{\beta_t} A_t^{-1/2} z_d, \|z_d\|_2 \leq 1 \},$$  \hspace{1cm} (17)

then plug in equation (15) to yield:

$$\hat{\theta}_t + \sqrt{\beta_t} A_t^{-1/2} z = \hat{\theta}_t + \sqrt{\beta_t} A_t^{-1/2} z - A_t^{-1} X_t^T b_t$$

$$= \hat{\theta}_t + \sqrt{\beta_t} A_t^{-1/2} z - A_t^{-1} (\sum_{s=1}^{t-1} b_s x_s).$$  \hspace{1cm} (18)

Therefore, we know that with high probability,

$$\theta^* \in \tilde{C}_t = \{ \hat{\theta}_t + \sqrt{\beta_t} A_t^{-1/2} z_d - A_t^{-1} (\sum_{s=1}^{t-1} b_s x_s) \}.  \hspace{1cm} (19)$$

Therefore, we have a computable confidence bound for $x$:

$$UCB_t(x) = \max_{\theta \in \tilde{C}_t} \langle x, \theta \rangle$$

$$= \langle x, \hat{\theta}_t \rangle + \max_{z_d \in B_d^2} \sqrt{\beta_t} \langle x, A_t^{-1/2} z_d \rangle - \sum_{s=1}^{t-1} b_s x^T A_t^{-1} x_s$$

$$\leq \langle x, \hat{\theta}_t \rangle + \sqrt{\beta_t} \|A_t^{-1/2} x\| - \sum_{s=1}^{t-1} b_s x^T A_t^{-1} x_s$$

$$\leq \langle x, \hat{\theta}_t \rangle + \sqrt{\beta_t} \|A_t^{-1/2} x\| + \epsilon \sum_{s=1}^{t-1} |x^T A_t^{-1} x_s|.$$  \hspace{1cm} (20)

The first inequality is derived by Cauchy Schwartz inequality and the fact that $z_d$ is in unit ball.  \hspace{1cm} \(\Box\)

### A.2.2 Proof of Lemma 10

**Proof.** Lemma 10 is a parallel to Theorem 5 in [Dani et al. (2008)](http://example.com), with the difference of sub-gaussian noise, ellipsoid centre $\hat{\theta}_t$ and misspecification in observation. The key idea is the same, namely to use induction to bound the growth of $Z_t = (\theta^* - \hat{\theta}_t)^T A_t (\theta^* - \hat{\theta}_t)$ and proof that $Z_t \leq \beta_t$, i.e. the $\theta^*$ is contained in $C_t$, at each time step $t$. The following analysis used the same notations and definitions as section 5.2 in [Dani et al. (2008)](http://example.com) unless otherwise specified. Under Lemma 10’s definition of confidence set $C_t$, we have that:

$$H_t = A_t (\hat{\theta}_t - \theta^*) = X_t^T \eta_t - \theta^*,$$  \hspace{1cm} (21)

$$Z_t = (\theta^* - \hat{\theta}_t)^T A_t (\theta^* - \hat{\theta}_t) = H_t^T A_t^{-1} H_t.$$  \hspace{1cm} (22)

Equation (21) holds because of this key property:

$$\hat{\theta}_t : A_t \hat{\theta}_t = X_t^T X_t \theta^* + X_t^T \eta_t.$$  \hspace{1cm} (23)

And the rest of the proof in [Dani et al. (2008)](http://example.com) should go through by substituting $Y_t$ with $H_t$ (defined above) and $\hat{\mu}$ with our definition of $\hat{\theta}$ (centre of the confidence ellipsoid). Except, to accommodate the sub-gaussian noise assumption that replaces their bounded noise assumption, we have to make two changes in the proof. Both are in analyzing the growth of $Z_t$ in the induction. Recall that [Dani et al. (2008)](http://example.com) proved this relation:

$$Z_t \leq Z_1 + 2 \sum_{\tau=1}^{t-1} \frac{\eta^2_{\tau} (\hat{\theta}_\tau - \theta^*)^2}{1 + w^2_{\tau}} + \sum_{\tau=1}^{t-1} \frac{\eta^2_{\tau} w^2_{\tau}}{1 + w^2_{\tau}}.$$  \hspace{1cm} (24)
We first look at the concentration of the sum of martingale difference sequence that makes up $Z_t$: same with [Dani et al. (2008)], define $M_t = 2\eta_t x^T (\theta_t - \theta^*)$ where $w_t \triangleq \sqrt{x^T A^{-1} x_t}$. According to our assumption, the noise sequence is a sub-gaussian martingale difference sequence with parameter $\sigma^2$. Therefore, $M_t$ is a sub-gaussian martingale difference sequence. Specifically, we know that the square of sub-gaussian parameter is $4\sigma^2 (\frac{x^T (\theta_t - \theta^*)}{1 + w_t^2})^2$. By definitions we know that $M_t | \mathcal{H}_t$ is $(\nu^2_t = 4\sigma^2 (\frac{x^T (\theta_t - \theta^*)}{1 + w_t^2})^2$, $a_t = 0)$ sub-exponential (definition 2.7 in Wainwright (2019)) and therefore the sum $\sum_{t=1}^T M_t$ is also sub-exponential, with parameters $(\sqrt{\sum_{t=1}^T \nu^2_t}, a = \max_{t} a_t = 0)$ (Theorem 2.19 (1) in Wainwright (2019)). The following inequality is conditioned on the fact that from time $\tau = 1 \ldots t, \theta^*$ is contained in $C_\tau$ (by the induction).

$$\sum_{\tau=1}^t \nu^2_t = 4\sigma^2 \sum_{\tau=1}^t \left( \frac{x^T (\theta_\tau - \theta^*)}{1 + w^2_\tau} \right)^2$$

$$\leq 4\sigma^2 \sum_{\tau=1}^t \left( \frac{\sqrt{\beta_t} w_\tau}{1 + w^2_\tau} \right)^2$$

$$\leq 4\sigma^2 \sum_{\tau=1}^t \beta_t (\min(1/2, w_\tau))^2$$

$$\leq 4\sigma^2 \sum_{\tau=1}^t \beta_t \min(1/4, w^2_\tau)$$

$$\leq 4\sigma^2 \beta_t \sum_{\tau=1}^t \min(1, w^2_\tau)$$

$$\leq 4\sigma^2 \beta_t (2d \ln(1 + t\kappa^2/d)) \text{ See bound 28}$$

$$= 8\sigma^2 d\beta_t \ln(1 + t\kappa^2/d).$$

The proof for the first three inequalities is the same as Lemma 7 and section 5.2.1 in [Dani et al. (2008)]. Then we apply a Bernstein-type concentration bound for sub-exponential martingale difference sequence (Theorem 2.19 (2) in [Wainwright (2019)]). Plugging in the values of $a$ and $\sum_{\tau=1}^t \nu^2_t$, we have that

$$\mathbb{P}(\sum_{\tau=1}^{t-1} M_\tau \geq s) \leq 2 \exp\left( \frac{-s^2}{2 \sum_{\tau=1}^{t-1} \nu^2_t} \right)$$

$$\leq 2 \exp\left( \frac{-s^2}{16\sigma^2 d\beta_t \ln(1 + (t - 1)\kappa^2/d)} \right)$$

$$s = \frac{\delta}{2\kappa^2} \leq 2 \exp\left( \frac{-\beta_t}{64\sigma^2 d \ln(1 + (t - 1)\kappa^2/d)} \right)$$

$$\leq \frac{\delta}{2\kappa^2} \text{ (Needed for union bound over all times).} \tag{25}$$

Therefore, as long as $\beta_t$ is larger or equal to $64\sigma^2 d \ln(1 + (t - 1)\kappa^2/d) \ln(4\kappa^2/\nu^2) \sum_{\tau=1}^{t-1} M_\tau \leq \frac{\beta_t}{2}$ with probability larger or equal to $1 - \frac{\delta}{2\kappa^2}$.

The second change is for the third quantity that makes up $Z_t$: $\sum_{\tau=1}^{t-1} \nu^2_t \frac{w^2_\tau}{1 + w^2_\tau}$. We need to bound $\max_{\tau \leq t-1} \eta^2_\tau$ with high probability. By algebra calculations, we know that $\eta^2_\tau$ is sub-exponential with parameters $(\nu = 32\sigma^4, a = 4\sigma^2)$ \textsuperscript{15} We can apply union bound with the tail bound of sub-exponential

\textsuperscript{15} For this part, we borrowed the proof from Example 2.8 in [Wainwright (2019)] and [http://proceedings.mlr.press/v33/honorio14-suppl.pdf]
variables:

\[ \mathbb{P}( \max_{\tau \leq t-1} (\eta_\tau^2 - \mathbb{E}[\eta^2]) \geq z) \leq \sum_{\tau=1}^{t-1} \mathbb{P}(\eta_\tau^2 - \mathbb{E}[\eta^2]) \geq z) \]

\[ \leq (t - 1) \exp\left( - \frac{z^2}{2d} \right) \quad (\text{Proposition 2.9 in Wainwright (2019)}) \]

\[ \leq \frac{\delta}{2t^2} \quad \text{(Needed for union bound over all times).} \]

Set \( z = 8\sigma^2 \ln(\frac{2t}{\delta}) \) so that \( \mathbb{P}(\max_{\tau \leq t-1} \eta_\tau^2 - \mathbb{E}[\eta^2] \leq z) = \mathbb{P}(\max_{\tau \leq t-1} \eta_\tau^2 \leq z + \mathbb{E}[\eta^2]) \geq 1 - \frac{\delta}{t^2} \).

By the fact that \( \mathbb{E}[\eta] = 0, \mathbb{E}[\eta^2] = \text{Var}(\eta) \leq \sigma^2 \), which is a property of subgaussian variables. So \( \mathbb{P}(\max_{\tau \leq t-1} \eta_\tau^2 \leq z + \sigma^2) \geq 1 - \frac{\delta}{t^2} \). The following holds with probability larger than \( 1 - \frac{\delta}{t^2} \):

\[ \sum_{\tau=1}^{t-1} \eta_\tau^2 \leq \frac{\max_{\tau \leq t-1} (\eta_\tau^2)}{\max_{\tau \leq t-1} (\eta_\tau^2) 2d \ln(1 + t\kappa^2/d)} \]

\[ = (8\sigma^2 \ln(\frac{2\kappa^3}{\delta}) + \sigma^2)2d \ln(1 + (t - 1)\kappa^2/d) \]

\[ = 8\sigma^2 \ln(\frac{2\kappa^3}{\delta}) + \frac{1}{8})2d \ln(1 + (t - 1)\kappa^2/d) \]

\[ = 16\sigma^2 d \ln(1 + (t - 1)\kappa^2/d) \left( \ln(\frac{2\kappa^3}{\delta}) + \frac{1}{8} \right). \]

Except the two changes above, one last thing to note is the quantity \( Z_1 \) analyzed at the end of proof of Lemma 12 in Dani et al. (2008). In our assumption of the reward function value, we conclude that

\[ Z_1 = (\theta^* - 0)^T I(\theta^* - 0) = \|\theta^*\|^2 \]

\[ = \sum_{i=1}^{d} (e_i^T \theta^*)^2 \quad (e_i \text{ is base vector of dimension } i, \text{ note that } e_i \in \mathcal{X}) \]

\[ \leq d(1 + \epsilon)^2. \]

As a result, if it is satisfied that \( Z_t \leq Z_1 + \beta_t/2 + 16\sigma^2 d \ln(1 + (t - 1)\kappa^2/d)ln(\frac{2\kappa^3}{\delta}) + \frac{1}{8} \) \leq \beta_t \), which enables the induction in Lemma 14 in Dani et al. (2008), then the rest of the proof should go through smoothly. We argue that setting \( \beta_t = C\sigma^2 d \ln(t) \ln(\frac{2\kappa^3}{\delta}) \) for a large enough constant \( C \) suffices. This is under the reasonable assumption that \( \epsilon = O(1) \) and \( \sigma = \text{a constant}\).

It is worth mentioning\(^{16}\) that Dani et al. (2008) requires the relationship between \( t \) and \( \delta \) to be approximately \( 0 < 1.05\delta \leq t^2 \), hence their requirement\(^{17}\) of “for sufficiently large \( T \)” in Theorem 1 and 2. This is because of the last step of their induction proof for Theorem 5 requires: \( Z_t \leq d + \beta_t/2 + 2d \ln(t) \leq \beta_t \). In our setting, the requirement in induction translates to this (second) constraint(plugging in \( \kappa^2 = d \)): \( \beta_t \geq 2d(1 + \epsilon)^2 + 32\sigma^2 d \ln(t)(\ln(\frac{2\kappa^3}{\delta}) + \frac{1}{8}) \).

Recall the first constraint on \( \beta_t \) is \( \beta_t \geq 64\sigma^2 d \ln(t)(\ln(\frac{4\kappa^3}{\delta})) \), from bound (25). Therefore, \( C \) should first satisfy \( C \geq 64 \) and for the second constraint we need\(^{19}\) \( C \geq \frac{3(1+\epsilon)^2}{4(\ln(2))^2}\sigma^2 + \frac{3}{\ln(2)} + 48 \). Therefore, the lower bound of \( C \) should depend on values of \( \epsilon \) and \( \sigma^2 \). The choice of \( C = 128 \) in the main theorem is an example that requires approximately \( \frac{1+\epsilon}{\sigma^2} \leq 7 \).

\(^{16}\)Recall that according to Lemma 2, \( \epsilon \) is bounded by the Lipschitz constant \( L \) and is therefore \( O(1) \).

\(^{17}\)This remark is made by Abbasi-Yadkori et al. (2011).

\(^{19}\)However, we believe that this should not translate to a constraint on \( t \), but on \( \delta \) instead. Because \( Z_t \leq \beta_t \) is required for every step \( t \) to complete the induction, so if it only holds for large \( t \) then the induction will fail as well.

\(^{19}\)This is from the second constraint: \( C\sigma^2 d \ln(t)(\ln(\frac{4\kappa^3}{\delta})) \geq \frac{3}{\ln(2)} + 48 \).
A.3 Proof of Theorem 4

Let us treat the number of bins/local algorithms $n$ as the input parameter to the algorithm. The regret bound of UCB-Meta (equation 4) should be independent of the input dimension $d$, given the dimension of the linear model $d(n)$. Therefore, throughout this proof we will abuse the notations and let $d$ denote the linear model dimension for simplicity.

**Proof.** First, we define the “good event” $E_{good}$ as an event where all confidence bound holds for all bins at all times. For a fixed bin, if $\mathbb{P}(\theta^{*} \notin \hat{C}_{t}, \exists t) \leq \delta/n$, as set in the algorithm, where $\hat{C}_{t} = \{\hat{\theta}_{t} + \sqrt{\frac{1}{t}A_{t}^{-1}z_{d} - A_{t}^{-1}(\sum_{s=1}^{t-1}b_{s}x_{s})}\}$ (Theorem 9), then by union bound, $\mathbb{P}(\theta^{*} \notin \hat{C}_{k,t}, \exists k) \leq \delta$, where $\hat{C}_{k,t}$ is the confidence ellipsoid of bin $k$ at time $t$. The good event is $E_{good} = \{\forall t, \forall k \in [n], \theta^{*}_{k,t} \in \hat{C}_{k,t}\}$. It happens with probability $\mathbb{P}(E_{good}) \geq 1 - \delta$, and the following proof will condition on it.

Here are some useful notations that make the proof easier to read: let $N^{k}(t)$ denote the number of times base-algorithm $A^{local}_{k}$ has been selected by (including) time $t$; let $k(t)$ denote the bin selected at time $t$; let $x_{t}$ denote the action selected at time $t$; let $\{\beta_{k,\cdot}\}, \{\alpha_{k,\cdot}\}$ and $\{\hat{\beta}_{k,\cdot}\}$ denote the set of parameters kept by that base-algorithm $A^{local}_{k}$. The upper confidence bound on value of the local linear function achieved by sub-algorithms at round $t$ is defined as $UCB_{k,(t),t}(x) = \langle x, \hat{\theta}_{k,N^{k}(t)} \rangle + \sqrt{\beta_{k,N^{k}(t)}}\|A_{k,N^{k}(t)}^{-1/2}x\| + \epsilon \sum_{\tau=1}^{N^{k}(t)-1} |x^{T}A_{N^{k}(t)}^{-1}x_{\tau}|$ for any action $x \in B_{k}$. Using the proof of Theorem 3, the good event hence indicates that for the base-algorithm selected at time $t$ and any action $x \in B_{k(t)}$:

$$UCB_{k,(t),t}(x) = \sqrt{\beta_{k,N^{k}(t)}}\|A_{k,N^{k}(t)}^{-1/2}x\| - 2\epsilon \sum_{\tau=1}^{N^{k}(t)-1} |x^{T}A_{N^{k}(t)}^{-1}x_{\tau}| - \langle x, \theta^{*}_{k} \rangle \leq UCB_{k,(t),t}(x).$$

By Lemma 2, the expected local function value $f(x)$ is bounded by

$$UCB_{k,(t),t}(x) = \sqrt{\beta_{k,N^{k}(t)}}\|A_{k,N^{k}(t)}^{-1/2}x\| - 2\epsilon \sum_{\tau=1}^{N^{k}(t)-1} |x^{T}A_{N^{k}(t)}^{-1}x_{\tau}| - \epsilon \leq f(x) \leq UCB_{k,(t),t}(x) + \epsilon.$$

A common way to bound pseudo regret for stochastic bandit is via Wald’s equality: $R_{T} = \sum_{n=1}^{N} \Delta_{n} \mathbb{E}[\tau_{n}(T)]$ where $\Delta_{n}(T)$ is the number of times arm $k$ gets pulled until time $T$, and $\Delta_{n}$ is the reward gap. We cannot trivially follow this, because the rewards of each bins are no longer i.i.d. Instead, we use this gap-independent decomposition for each bin $k$:

$$R_{k} = \sum_{t \text{ bin } = k} \left( f^{*} - f_{x_{t}=B_{k}(x_{t})} \right)$$

$$= \sum_{t \text{ bin } = k} \left( f^{*} - UCB_{A_{k},t} + UCB_{A_{k},t} - f(x_{t}) \right)$$

$$= \sum_{t \text{ bin } = k} \left( f^{*} - UCB_{A_{k},t} + UCB_{k,t} + \epsilon - f(x_{t}) \right)$$

$$\leq \sum_{t \text{ bin } = k} \left( UCB_{k,t} + \epsilon - f(x_{t}) \right)$$

$$\leq \sum_{t \text{ bin } = k} \left( \sqrt{\beta_{k,N^{k}(t)}}\|A_{k,N^{k}(t)}^{-1/2}x_{t}\| + \epsilon \sum_{\tau=1}^{N^{k}(t)-1} |x_{t}^{T}A_{N^{k}(t)}^{-1}x_{\tau}| + 2\epsilon \right)$$

$$= \sum_{s=1}^{N^{k}(T)} \left( \sqrt{\beta_{s,N^{k}(t)}}\|A_{s,N^{k}(t)}^{-1/2}x_{t}\| + 2\epsilon \sum_{\tau=1}^{s-1} |x_{t}^{T}A_{s}^{-1}x_{\tau}| + 2\epsilon \right).$$

The first inequality holds because of the algorithm’s bin selection rule: if bin $B_{k}$ is chosen then $f^{*} \leq UCB_{k+1,t} \leq UCB_{k,t}$ . By the bounded function value assumption, $f^{*} - f_{x_{t}=B_{k}(x_{t})} \leq 2$, ...
where we used Lemma 11. Now we can bound term \#1 using bound (28).

**Proof of Lemma 11.**

First we establish this bound the same way as Dani et al. (2008). Namely, for any local misspecified linear bandit algorithm that is ran \(T\) times with data \((x_t, y_t)_{t=1,...,T}\),

\[
\sum_{t=1}^{T} \| x_t^T A_t^{-1} x_t \| \wedge 1 \leq 2 \ln \left( \prod_{t=1}^{T} (1 + x_t^T A_t^{-1} x_t) \right) = 2 \ln \left( \prod_{t=1}^{T} \frac{\det(A_{t+1})}{\det(A_t)} \right) = 2 \ln \left( \frac{\det A_{T+1}}{\det A_1} \right) \leq 2 \ln (1 + (T\kappa^2/d)^d) = 2d \ln (1 + T\kappa^2/d),
\]

where we used Lemma [11]. Now we can bound term \#1 using bound (28).

\[
\sum_{s=1}^{N^k(T)} 2(\sqrt{\beta_{k,s}} \| A_{k,s}^{-1/2} x_t \| \wedge 1) \leq \sqrt{N^k(T)} \sum_{s=1}^{N^k(T)} 4(\beta_{k,s} \| x_t^T A_{k,s}^{-1} x_{k,s} \| \wedge 1) \leq \frac{4\beta_{k,N^k(T)}N^k(T) \sum_{s=1}^{N^k(T)} \| x_t^T A_{k,s}^{-1} x_{k,s} \| \wedge 1}{4\beta_{k,N^k(T)}N^k(T)2\ln \left( \prod_{s=1}^{N^k(T)} (1 + x_t^T A_{k,s}^{-1} x_{k,s}) \right)} = \frac{4\beta_{k,N^k(T)}N^k(T)2\ln \left( \frac{\det(A_{N^k(T)+1})}{\det(A_1)} \right)}{8d\beta_{k,N^k(T)}N^k(T) \ln (1 + N^k(T)\kappa^2/d) = \sqrt{8d\beta_{k,N^k(T)}N^k(T) \ln (1 + N^k(T))}.}
\]

**Lemma 11.** For \(t \geq 1, 1 + x_t^T A_t^{-1} x_t = \det(A_{t+1})/\det(A_t).\) Also, \(\det(A_t) \leq (1 + (t-1)\kappa^2/d)^d.\)

**Proof of Lemma 11**

\[
\det(A_{t+1}) = \det(A_t(I_d + A_t^{-1} x_t x_t^T)) = \det(A_t) \det(I_d + A_t^{-1} x_t x_t^T) = \det(A_t) \det(I_1 + x_t^T A_t^{-1} x_t) = \det(A_t)(1 + x_t^T A_t^{-1} x_t).
\]
The third equation uses Sylvester’s determinant theorem: 
\[ \det(I_m + A_{m \times n} B_{n \times m}) = \det(I_n + B_{n \times m} A_{m \times n}). \]
The trace of a matrix is the product of its eigenvalues and the determinant is the sum of eigenvalues, and for the trace of the positive definite matrix \( A_t \) we have,
\[ \text{tr}(A_t) = \text{tr}(I + \sum_{\tau} x_\tau x_\tau^T) = d + \sum_{\tau} \|x_\tau\|_2^2 \leq d + (t - 1)\kappa^2. \]
Therefore, using the inequality of arithmetic and geometric mean, 
\[ \det(A_t) \leq (1 + (t - 1)\kappa^2/d)^d. \]

Summing over all the suboptimal bins, we have that
\begin{align*}
\sum_{k=1}^{n-1} \sum_{s=1}^{N^k(T)} 2(\sqrt{\beta_{k,s}} A_{k,s}^{-1/2} x_{k,s} \wedge 1) &\leq \sum_{k=1}^{n} 8d\beta_{k,N^k(T)} N^k(T) \ln (1 + N^k(T)) \\
&\leq \sqrt{\sum_{k=1}^{n} N^k(T) \sum_{k=1}^{n} 8d\beta_{k,N^k(T)} \ln (1 + N^k(T))} \\
&= \sqrt{T \sum_{k=1}^{n} 8d\beta_{k,N^k(T)} \ln (1 + N^k(T))} \\
&\leq \sqrt{8d T \ln (1 + T)}. \quad (29)
\end{align*}

**A.3.2 High probability regret bound part II (term #2)**

Here we directly call previous result in bound (14), but replace the total number of step with \( N^k(T) \), the number of pulls for one fixed bin \( k \). We have for term #2,
\[ \sum_{s=1}^{N^k(T)} \sum_{\tau=1}^{t-1} 2\varepsilon \|x_{k,s}^{-1} x_{k,\tau}\| \leq 2\varepsilon N^k(T) d \sqrt{2 \ln(1 + N^k(T))}. \]

Summing over all suboptimal bins, we have that
\begin{align*}
\sum_{k=1}^{n} 2\varepsilon N^k(T) d \sqrt{2 \ln(1 + N^k(T))} &\leq 2\varepsilon d \sqrt{2 \ln(1 + T)} \sum_{k=1}^{n} N^k(T) \\
&= 2\varepsilon d T \sqrt{2 \ln(1 + T)}. \quad (30)
\end{align*}

**A.3.3 Putting it together**

Combining the decomposition in equation (27) and the results in subsections A.3.1 and A.3.2, we have a high probability regret bound for the UCB-Meta-algorithm:
\[ R_T = \sum_{k=1}^{n} R_k \leq \sqrt{8d T n \beta_T \ln (1 + T^2) + 2\varepsilon d T \sqrt{2 \ln(1 + T^2)}) + 2\varepsilon T} = \mathcal{O}(d \ln(T) \sqrt{T n \ln(T^2 n / \delta)} + \varepsilon T \sqrt{\ln(T) + \varepsilon T}). \quad (31)\]
The last step plugs in \( \beta_T = \mathcal{O}(d \ln(T) \ln(T^2 n / \delta)). \)
A.3.4 Proof of Theorem 5

Proof. Algorithm 3 executes Algorithm 2 for a sequence of pre-defined time periods, \( \{T_i = 2^i, i = 0, 1, \ldots N\} \). At the beginning of each period, the update history is cleared and the number of arms \( n \) is reset with respect to the current horizon \( T_i \). However, since we would like to acquire a high-probability regret bound after applying the doubling trick, we need to set the fail probability of Meta-algorithms during period \( i \) to \( \delta_i = 6\delta/\pi^2 i^2 \). Using a union bound, we can conclude the following (\( R_i(T_i) \) denotes the regret incurred in time period \( i \) of length \( T_i \) only).

\[
P(\forall i, \text{the bound hold for } R_i(T_i)) = 1 - \sum_i P(\text{the bound does not hold for } R_i(T_i)) = 1 - \sum_i \frac{6\delta}{\pi^2 i^2} \approx 1 - \delta.
\]

In the last step we use the fact that the sum of sequence \( \sum_i^\infty \frac{1}{i^2} \) converges to \( \pi^2/6 \).

Now, the total regret is simply a summation over \( i \). The following holds with probability \( 1 - \delta \),

\[
R(T) \leq \sum_{i=1}^N R_i(T_i)
\leq \sum_{i=1}^N \tilde{O}(dT_i^\alpha) = \tilde{O} \left( d \sum_{i=1}^N 2^{i\alpha} \right)
\leq \tilde{O} \left( d2^{\alpha(N-1)} \right)
= \tilde{O}(dT^{\alpha}). \quad (32)
\]

At step 4, the number of time periods \( N \) is the smallest integer such that \( \sum_{i=0}^N 2^i \geq T \), so \( N = 1 + \lceil \log_2(T) \rceil \). The sum of geometric sequence is \( 2^{\alpha \lceil \log_2(T) \rceil} = (2^{\alpha \log_2(T) + \alpha})^a = T^a \) for some constant \( c \) smaller than 1. Also, note that step 2 holds even though the fail probability is changed to \( \delta_i = 6\delta/\pi^2 i^2 \) is because as specified in Theorem 4, the term \( \delta \) appears in a log term and the maximum value of \( 1/\delta_N = \pi^2 \log_2(T)/6\delta \), therefore the extra factor caused by smaller \( \delta \) to the regret is still a log term of \( T_i \) and omitted in the proof here.

Bound (32) suffices to say that meta-algorithm with doubling trick has the same regret rate as meta-algorithm with known horizon, with some additional constant factors suffered from restarting. \( \square \)

A.4 Proof of Theorem 6

Proof. Here we prove that Corral with smooth-wrapper is applicable to this task and achieves minimax expected regret rate apart from log factors. We directly use the proof of Theorem 5.3 in Pacchiano et al. (2020) and their notations. \( \delta \) is the fail probability, \( M \) is the number of base-algorithms, \( \rho \) is the reciprocal of the smallest possibility for base-algorithms over the \( T \) rounds and \( \eta \) is the learning rate. \( U(T, \delta) \) is the high probability bound of the selected base-algorithm. The regret of Corral with smooth wrapper is bounded by:

\[
R(T) \leq \tilde{O} \left( \frac{M \ln(T)}{\eta} + T\eta \right) + \delta T + 8\sqrt{MT \log \left( \frac{4TM}{\delta} \right)} - \mathbb{E} \left[ \frac{\rho}{40\eta \ln(T)} - 2\rho U(T/\rho, \delta) \log(T) \right],
\]

and we know from Theorem 3 in our paper that the base algorithm (Algorithm 1) that locates in the global maximum’s bin has anytime high probability regret bound \( \tilde{U}(T, \delta) = \tilde{O}(cT d(\alpha) + \)
We minimize this by setting the derivative w.r.t. \( \tilde{\text{smooth}} \). The bound of the Corral algorithm holds with respect to any of its base-algorithm with high probability regret difference between the rates achieved for \( \hat{\alpha} \) over a grid for the unknown point \( \hat{\alpha} \).

**Proof.** There exists an input parameter \( \alpha \), the algorithm sets \( \alpha \), \( \alpha \) is also \( \alpha \). If the input parameter \( \alpha \) is smaller than \( \alpha \), then the calculated misspecification error \( \epsilon \) is smaller than the true \( \epsilon^* = \hat{\alpha}(T \frac{\rho}{\log(T)}) \), causing the confidence bound to be invalid. Therefore, the regret bound does not hold for when \( \alpha' > \alpha \). If the input parameter is smaller than \( \alpha \), then we can simply use the fact that functions that are \( \alpha \)-Hölder smooth are also \( \alpha \)-Hölder smooth: \( H(\alpha, L) \subset H(\alpha', L) \). Therefore, the regret of the algorithm with input parameter \( \alpha' \leq \alpha \) is bounded by \( R(T) \leq \hat{\alpha}(d(\alpha') (\sqrt{Tn} + \epsilon T)) = \hat{\alpha}(d(\alpha') T \frac{d+4}{\rho^2 + \rho^4}) \).

### A.5 Proof of Lemma 7

**Proof.** According to Theorem 4, the algorithm sets \( \alpha = T \frac{\rho}{\log(T)} \frac{d+4}{\rho^2 + \rho^4} \), for any true \( \alpha \) in \( (0, R] \). There are two sources that made up the cost of adaptation when using Corral. The first one is the cost of searching over a grid for the unknown point \( \hat{\alpha} \). The second one is the cost of approximation, specifically the difference between the rates achieved for \( \hat{\alpha} \) and the true \( \alpha \). We will first derive the cost of grid search. As specified in the proof of Theorem 5.3 in Pacchiano et al. [2020], the following bound of regret of the Corral algorithm holds with respect to any of its base-algorithm with high probability regret bound \( U(T, \delta) \). The notations were introduced in Appendix section A.4.

\[
R(T) \leq \mathcal{O}(\frac{M \ln(T)}{\eta} + T\eta) - \mathbb{E}[-\frac{\rho}{40\eta \ln(T)} - 2\rho U(T / \rho, \delta) \log(T)] + \delta T + 8 \sqrt{MT \log(\frac{4TM}{\delta})}.
\]

(36)

Plugging the regret rate of base-algorithm in Lemma 7 the expected pseudo-regret of Corral with smooth wrapper is therefore bounded by:

\[
R(T) \leq \mathcal{O}(\frac{M \ln(T)}{\eta} + T\eta + \sqrt{MT}) + \delta T - \mathbb{E}[-\frac{\rho}{40\eta \ln(T)} - 2\rho \mathcal{O}(\frac{T}{\rho} \frac{d+4}{\rho^2 + \rho^4})] \log(T)]
\]

(37)
Similarly, we first maximize over \( \rho \) by setting the derivative w.r.t \( \rho \) to zero by setting \( \rho = \tilde{O}(\eta^{-\alpha} d^{\frac{d+2}{\alpha}}) \). Then the above rate comes to

\[
R(T) \leq \tilde{O}\left( \frac{M}{\eta} + T\eta + \sqrt{MT + d^{\frac{d+2}{\alpha}} T \eta^{\frac{\alpha}{\alpha + \sigma}}} \right).
\]

(38)

However, since \( \eta \) is a parameter of the Corral algorithm which does not know \( \hat{\alpha} \) or \( \alpha \), we will rely on the parameter \( R \) specified by the user. Let us set \( \eta \) with respect to \( \alpha = R \), i.e. \( \eta = \tilde{O}(d^{-1} T^{-\frac{d+2}{\alpha}}) \), and plug in the number of grid points (base-algorithms) \( M = \lceil \log(T) \rceil \).

\[
\tilde{O}\left( \frac{M}{\eta} + T\eta + \sqrt{MT + d^{\frac{d+2}{\alpha}} T \eta^{\frac{\alpha}{\alpha + \sigma}}} \right) = \tilde{O}(dT^{\frac{d+2}{\alpha + \sigma} + d^{2+2Rd+R\alpha}}).
\]

(39)

It is obvious that this rate is not the minimax optimal rate for class \( \sum(\hat{\alpha}) \), this gap shows the cost of grid search.

Next, let us consider the cost of approximation and how it is eliminated by using the linear grid \((\text{Hoffmann et al., 2011})\). Namely, we show that adaptation for \( \hat{\alpha} \) is equivalent to adaptation for \( \alpha \):

\[
\tilde{O}(dT^{\frac{d+2}{\alpha + \sigma} + d^{\frac{d+2}{\alpha}}}) = \tilde{O}(dT^{\frac{d+2}{\alpha + \sigma} + d^{\frac{d+2}{\alpha}} + d^{2+2Rd+R\alpha}}).
\]

(40)

The equality holds because \( |\alpha - \hat{\alpha}| \leq \frac{R}{\log(T)} \). Let \( J = d^{\frac{d+2}{2Rd+R\alpha}} + \frac{R}{d+\alpha} \) and \( Q = d^{\frac{d+2}{2Rd+R\alpha}} + \frac{R}{d+\alpha} \), then \( W \triangleq \frac{T - J}{T} \leq T \frac{d+\alpha}{d+2Rd+R\alpha} \). Taking the log of \( W \) yields \( \log(W) = R \frac{d+2}{(d+2Rd+R\alpha)(d+\alpha)} \). Since both \( \alpha \) and \( \hat{\alpha} \) are bounded by a constant range \((0, 2]\), the term \( d^{\frac{d+2}{2Rd+R\alpha}} \frac{R}{d+\alpha} \) is bounded by a constant range \((0, 2]\), the term \( R \frac{d+2}{(d+2Rd+R\alpha)(d+\alpha)} \) is bounded by a constant range \((0, 2]\).

Therefore, for functions with Hölder exponent \( \alpha < R \), the second term in equation (40) is the dominant term and the expected regret rate is \( \tilde{O}(dT^{\frac{d+2}{\alpha + \sigma} + d^{\frac{d+2}{\alpha}}}) \). For functions with Hölder exponent \( \alpha \geq R \), which essentially belongs to a subset of \( \sum(R, L) \), they will all have the same rate which is \( \tilde{O}(dT^{\frac{d+2}{\alpha + \sigma}}) \). When \( \alpha = R \), this matches the minimax rate for \( \alpha \). \( \square \)
B  Additional algorithms for the main document

B.1 Doubling procedure for Algorithm 2

**Algorithm 3** Doubling procedure for Algorithm 2

**Require:** Meta-algorithm $A_{global}$ (Algorithm 2), fail probability $\delta$

1: for $i = 0 \ldots$ do
2: $T_i = 2^i$
3: Restart $A_{global}$ with initialization parameters $n_i = \lceil T_i^{\frac{\alpha}{\pi^2 \delta}} / \ln(T_i) \rceil$ and fail probability $\delta_i = 6\delta / \pi^2 i^2$
4: Run $A_{global}$ for $T_i$ steps.
5: end for

B.2 The Corral Master algorithm

For easier reference, we include the copy of Algorithm 7 in Pacchiano et al. (2020).

**Algorithm 4** Corral Master (Algorithm 7 in Pacchiano et al. (2020))

**Require:** Base algorithms $\{B_j\}_{j=1}^M$, learning rate $\eta$.

1: Initialize: $\gamma = 1/T, \beta = e^{\frac{1}{\eta T}}$, $\eta_t, j = \eta, \rho_1^t = 2M, p_1^t = \frac{1}{\rho_1^t}, p_1^t = 1/M$ for all $j \in [M]$.
2: for $t = 1, \ldots, T$ do
3: Sample $i_t \sim p_t$.
4: Receive feedback $r_t$ from base $B_{i_t}$.
5: Update $p_t, \eta_t$ and $p_{t+1}, \eta_{t+1}$ and $p_{t+1}$ using $r_t$ via Corral-Update (takes input $\eta_t, p_t, \beta$, lower bound $p_t$ and current feedback $r_t$).
6: for $j = 1, \ldots, M$ do
7: Set $\rho_{t+1}^j = \frac{1}{p_{t+1}^j}$.
8: end for
9: end for

The corral update procedure is in Algorithm 5 and the smooth wrapper for the base-algorithms in Algorithm 3 in Pacchiano et al. (2020).