Ordinal Embedding Relaxations
Parameterized Above Tight Lower Bound✩

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Abstract
We study ordinal embedding relaxations in the realm of parameterized complexity. We prove the existence of a quadratic kernel for the BETWEENNESS problem parameterized above tight lower bound, which is stated as follows. For a set $V$ of variables and set $C$ of constraints “$v_i$ is between $v_j$ and $v_k$”, decide whether there is a bijection from $V$ to the set $\{1, \ldots, |V|\}$ satisfying at least $|C|/3 + \kappa$ of the constraints in $C$. Our result solves an open problem attributed to Benny Chor in Niedermeier's monograph “Invitation to Fixed-Parameter Algorithms.” An approach developed in this paper can be used to determine parameterized complexity of a number of other optimization problems on permutations parameterized above or below tight bounds.

Key words: Fixed-parameter Tractability, Parameterized Complexity, Kernel, Probabilistic Method

1. Introduction

The problem of mapping points with measured pairwise distances into a target metric space has a long history and been studied extensively from multiple perspectives due to its numerous applications. The quality of such an embedding can be measured with various objectives; for example isometric embeddings preserve all distances while aiming at low-dimensional target spaces. Yet, for many contexts in nearest-neighbor search, visualization, clustering and compression it is the order of distances rather than the distances themselves that captures the relevant information. The study of such ordinal embeddings dates back to the 1950’s and has recently witnessed a surge in interest [1, 2, 3].

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In an ordinal embedding the relative order between pairs of distances must be preserved as much as possible, i.e., one minimizes the relaxation of an ordinal embedding defined as the maximum ratio between two distances whose relative order is inverted by the embedding.

Here we study the one-dimensional ordinal embedding of partial orders that specify the maximum edge for some triangles. This problem has been studied under the name of Betweenness (see the first problem in A12 of [14]), which takes a set $V$ of variables and a set $C$ of betweenness constraints of the form “$v_i$ is between $v_j$ and $v_k$” for distinct variables $v_i, v_j, v_k \in V$. Such a constraint will be written as $(v_i, \{v_j, v_k\})$. The objective is to find a bijection $\alpha$ from $V$ to the set $\{1, \ldots, |V|\}$ that “satisfies” the maximum number of constraints from $C$, where a constraint $(v_i, \{v_j, v_k\})$ is satisfied by $\alpha$ if either $\alpha(v_j) < \alpha(v_i) < \alpha(v_k)$ or $\alpha(v_k) < \alpha(v_i) < \alpha(v_j)$ holds. We also refer to $\alpha$ as a linear arrangement of $V$.

Such linear arrangements are of significant interest in molecular biology, where for example markers on a chromosome need to be totally ordered as to satisfy the maximum number of constraints [11, 15]. More theoretical interest comes from the constraint programming framework with unbounded domains and interval graph recognition [21].

Despite its simple formulation, the Betweenness problem has a challenging computational complexity. Already deciding if all constraints can be satisfied by some linear arrangement is an NP-complete problem [26]. Hence, the maximization problem is NP-hard. On the other hand, any uniformly random permutation of the variables satisfies at least one-third of all constraints, and this fraction is tight. Better approximation ratios are hard to achieve: the fraction of one-third is best-possible under the Unique Games Conjecture [9], and it is NP-hard to find a linear arrangement that satisfies a $1 - \epsilon$ fraction of the constraints for some $\epsilon \in (0, 1/48)$ [10]. The mere positive result is a polynomial time algorithm that either determines that there is no linear arrangement that satisfies all $m$ constraints or finds a linear arrangement satisfying at least half of them [10, 24].

A parameterized problem is a subset $L \subseteq \Sigma^* \times \mathbb{N}$ over a finite alphabet $\Sigma$. $L$ is fixed-parameter tractable if the membership of an instance $(x, \kappa)$ in $\Sigma^* \times \mathbb{N}$ can be decided in time $|x|^{O(1)} \cdot f(\kappa)$ where $f$ is a computable function of the parameter $|x| \leq \kappa$. If the nonparameterized version of $L$ (where $\kappa$ is just a part of input) is NP-hard, then the function $f(\kappa)$ must be superpolynomial provided $P \neq \text{NP}$. Often $f(\kappa)$ is “moderately exponential,” which makes the problem practically feasible for small values of $\kappa$. Thus, it is important to parameterize a problem in such a way that the instances with small values of $\kappa$ are of real interest.

One can parameterize the Betweenness problem in the standard way by asking whether there exists a linear arrangement that satisfies at least $\kappa$ of the constraints, where $\kappa$ is the parameter. This parameterized problem is trivially fixed-parameter tractable for the following simple reason. A uniformly random permutation of the variables in $V$ satisfies one-third of the constraints in expectation; thus if $\kappa \leq |C|/3$ then the answer is “yes” whereas if $\kappa > |C|/3$ then the
instance size is bounded by a function of $\kappa$ and we can solve the problem by brute-force. Note that the above standard parameterization is of little value as the parameter $\kappa$ will often be large. Thus, it makes sense to consider the following natural parameterization above a tight lower bound of the Betweenness problem [25].

Observe that $|C|/3$ is a lower bound on any optimal solution. On the other hand, for a set $C$ of constraints containing all three possible constraints on each 3-set of variables, no more than $|C|/3$ of the constraints in $C$ can be satisfied in any linear arrangement. Hence the lower bound of one-third on the fraction of satisfiable constraints is tight, in the sense that it is attained by an infinite family of instances. So the right question to ask is whether there exists a linear arrangement that satisfies at least $|C|/3 + \kappa$ of the constraints. The parameterized complexity of this problem attributed to Benny Chor is open, and was stated as such in [25]. Since the Betweenness problem is $\text{NP}$-complete, the complementary question of whether all but $\kappa$ constraints are satisfiable by some linear arrangement is not fixed-parameter tractable, unless $\text{P} = \text{NP}$. For the special case of a dense set of constraints, containing a constraint for each 3-subset of variables, a subexponential fixed-parameter algorithm was recently obtained [4].

Parameterized complexity of maximization problems parameterized above tight lower bounds often tends to be quite a challenging question, with only a handful of them answered. A prominent example is to decide whether a planar graph $G$ contains an independent set of size $|V(G)|/4 + \kappa$; it is unknown whether this problem is fixed-parameter tractable or not. Further addressed were above-tight-lower-bound parameterizations of $\text{MAX-CUT}$, $\text{MAX-SAT}$, $\text{MAX-r-SAT}$, and $\text{LINEAR ORDERING}$ [22, 2, 18]. Interestingly, all these problems were shown to be fixed-parameter tractable. Many open problems in that area are listed in [23]: in particular, the parameterized complexity of $\text{MAX LIN-2}$ parameterized above a tight lower bound remains unknown despite an effort in [18].

In this paper we settle Benny Chor’s question [25, p. 43] about the parameterized complexity of the following problem:

**Betweenness Above Tight Lower Bound (BATLB)**

**Instance**: a set $C$ of betweenness constraints over variables $V$ and an integer $\kappa \geq 0$.

**Parameter**: The integer $\kappa$.

**Question**: Is there a bijection $\alpha : V \rightarrow \{1, \ldots, |V|\}$ that satisfies at least $|C|/3 + \kappa$ constraints from $C$, that is, for at least $|C|/3 + \kappa$ constraints $(v_i, \{v_j, v_k\}) \in C$ we have either $\alpha(v_j) < \alpha(v_i) < \alpha(v_k)$ or $\alpha(v_k) < \alpha(v_i) < \alpha(v_j)$?

Our main result is that BATLB is fixed-parameter tractable. Moreover, we show that BATLB has a problem kernel of quadratic size, namely, any instance is polynomial-time reducible to an equivalent instance of size $O(\kappa^2)$. (We give a formal definition of a problem kernel in the next paragraph.)
is obtained via a nontrivial extension of the recently introduced probabilistic
Strictly Above/Below Expectation Method (SABEM) \[18\], which shows fixed-
parameter tractability of Linear Ordering and three special cases of Max Lin-2 parameterized above tight lower bounds. After the first version of this paper was completed (that version was not much different from the current one, see \url{http://arxiv.org/abs/0907.5427v1}), another paper \[2\] that applies SABEM was written. Alon et al. \[2\] further developed SABEM to prove fixed-
parameter tractability of Max \(r\)-SAT parameterized above a tight lower bound, but a simple modification of SABEM in \[2\] is not applicable to BATLB, see Section 2.

We describe SABEM briefly in Section 2 and point out how to extend it in
order to obtain a quadratic kernel for BATLB. The necessity to extend SABEM
lies with the fact that a feasible solution to BATLB is a permutation of vari-
ables; see Section 2 for a “high-level” discussion and Section 3 for details. Note
that our extension of SABEM can be used to determine parameterized complexity
of a number of other optimization problems on permutations parameterized
above or below tight bounds, cf. \[19\].

Given a parameterized problem \(L\), a kernelization of \(L\) is a polynomial-time
algorithm that maps an instance \((x, \kappa)\) to an instance \((x', \kappa')\) (the (problem)
kernel) such that (i) \((x, \kappa)\) \(\in\) \(L\) if and only if \((x', \kappa')\) \(\in\) \(L\), (ii) \(\kappa' \leq f(\kappa)\), and
(iii) \(|x'| \leq g(\kappa)\) for some functions \(f\) and \(g\). The function \(g(\kappa)\) is called the
size of the kernel. A parameterized problem is fixed-parameter tractable if and
only if it is decidable and admits a kernelization \[13\]; however, the problem
kernels obtained by this general result have impractically large size. Therefore,
one tries to develop kernelizations that yield problem kernels of smaller size\[1\].
A survey of Guo and Niedermeier \[16\] on kernelization lists some problem for
which polynomial size kernels were obtained. However, polynomial size kernels
are known only for some fixed-parameter tractable problems and Bodlaender
et al. \[7\] proved that many fixed-parameter tractable problems do not have
polynomial size kernels under reasonable complexity-theoretic assumptions.

2. Strictly Above/Below Expectation Method

The Strictly Above/Below Expectation method was recently introduced in
\[18\] to prove fixed-parameter tractability of maximization (minimization, re-
spectively) problems \(\Pi\) parameterized above (below, respectively) tight lower
(upper, respectively) bounds. In that method, we first apply some reductions
rules to reduce the given problem \(\Pi\) to its special case \(\Pi'\). Then we introduce a
random variable \(X\) such that the answer to \(\Pi'\) is Yes if and only if \(X\) takes with
positive probability a value greater or equal to the parameter \(\kappa\). If \(X\) happens to
be a symmetric random variable then the simple inequality \(\Pr(X \geq \sqrt{E[X^2]}) > 0\)

\[1\] Problem kernels are of great practical importance when it comes to solving \(\text{NP}\)-hard
problems; they can be interpreted as polynomial-time preprocessing with a quality guaran-
tee \[18, 17\].
can be useful; here $P(\cdot)$ and $E[\cdot]$ denote probability and expectation, respectively. An application is the Linear Ordering problem and a special case of Max Lin-2 parameterized above tight lower bounds [18]. However, often $X$ is not symmetric. In many such cases the following lemma by Alon, Gutin and Krivelevich [3] is of use; a weaker version of this lemma was obtained by Hästad and Venkatesh [20].

**Lemma 2.1.** Let $X$ be a real random variable and suppose that its first, second and forth moments satisfy $E[X] = 0$, $E[X^2] = \sigma^2 > 0$ and $E[X^4] \leq b\sigma^4$, respectively. Then $P(X > \frac{\sigma^4}{4} \sqrt{b}) \geq \frac{1}{4\sqrt{b}}$.

Often this result can be combined with the next lemma, which is an extension of Khinchin’s Inequality by Bourgain [8].

**Lemma 2.2.** Let $f = f(x_1, \ldots, x_n)$ be a polynomial of degree $r$ in $n$ variables $x_1, \ldots, x_n$ with domain $\{-1, 1\}$. Define a random variable $X$ by choosing a vector $(\epsilon_1, \ldots, \epsilon_n) \in \{-1, 1\}^n$ uniformly at random and setting $X = f(\epsilon_1, \ldots, \epsilon_n)$. Then, $E[X^4] \leq 2^{6r}(E[X^2])^2$.

These two lemmas were used in [2] and [18] to prove fixed-parameter tractability of Max $r$-SAT and a special case of Max Lin-2 parameterized above tight lower bounds. Unfortunately, it appears to be impossible to introduce a random variable $X$ for which $P(X \geq k) > 0$ if and only if the answer to BATLB is Yes and such that $X$ is either symmetric or satisfies the conditions of Lemma 2.2. Thus, in the next section, we introduce $X$ for which we have a weaker property with respect to BATLB: if $P(X \geq k) > 0$ then the answer to BATLB is Yes. This $X$, however, satisfies the conditions of Lemma 2.2.

To apply Lemma 2.1 we need to evaluate $E[X^2]$. While such evaluations in [2] and [18] are rather straightforward, our evaluation of $E[X^2]$ is quite involved and requires assistance of a computer. Note that we cannot algebraically simplify $X$ as was done in [2] in order to simplify computation of $E[X^2]$. This is due to the weaker property of $X$ with respect to BATLB: for some Yes-instances of BATLB we may have $P(X \geq k) = 0$.

### 3. A quadratic kernel for BATLB

We will now show fixed-parameter tractability of BATLB. In fact, we will prove a stronger statement, i.e., that this problem has a kernel of quadratic size.

For a constraint $C$ of $C$ let $\text{vars}(C)$ denote the set of variables in $C$. We call a triple $A, B, C$ of distinct betweenness constraints complete if $\text{vars}(A) = \text{vars}(B) = \text{vars}(C)$.

Consider the following reduction rule: if $C$ contains a complete triple of constraints, delete these constraints from $C$ and delete from $V$ any variable that appears only in the triple. Since for every linear arrangement exactly one constraint in each complete triple is satisfied we have the following:
Lemma 3.1. Let \((V, C)\) be an instance of \text{BATLB} and let \((V', C')\) be obtained from \((V, C)\) by applying the reduction rule as long as possible. Then \((V, C)\) is a Yes-instance of \text{BATLB} if and only if so is \((V', C')\).

An instance \((V, C)\) of \text{BATLB} is irreducible if it does not contain a complete triple. Observe that using Lemma 3.1 we can transform any instance into an irreducible one and it will take no more than \(O(m^3)\) time.

Consider an instance \((V, C)\), for a set \(V\) of variables and a set \(C = \{C_1, \ldots, C_m\}\) of betweenness constraints, and a random function \(\phi: V \rightarrow \{0, 1, 2, 3\}\). (The reason we consider a random function \(\phi: V \rightarrow \{0, 1, 2, 3\}\) rather than a random function \(\phi: V \rightarrow \{0, 1\}\) is given in the end of this section.) Let \(\ell_i(\phi)\) be the number of variables in \(V\) mapped by \(\phi\) to \(i\) for \(i = 0, 1, 2, 3\). Now obtain a bijection \(\alpha: V \rightarrow \{1, \ldots, |V|\}\) by randomly assigning values \(1, \ldots, \ell_0(\phi)\) to all \(\alpha(v)\) for which \(\phi(v) = 0\), and values \(\sum_{j=0}^{i-1} \ell_i(\phi) + 1, \ldots, \sum_{j=0}^{i} \ell_i(\phi)\) to all \(\alpha(v)\) for which \(\phi(v) = j\) for every \(j = 1, 2, 3\). We call such a linear arrangement \(\alpha\) a \(\phi\)-compatible bijection. It is easy to see that \(\alpha\) obtained in this two stage process is, in fact, a random linear arrangement, but this fact is not going to be used here.

Now assume that a function \(\phi: V \rightarrow \{0, 1, 2, 3\}\) is fixed and consider a constraint \(C_p = (v_i, \{v_j, v_k\}) \in C\). Let \(\alpha\) be a random \(\phi\)-compatible bijection and \(v_p(\alpha) = 1\) if \(C_p\) is satisfied and 0, otherwise. Let the weights \(w(C_p, \phi) = \mathbb{E}(v_p(\alpha)) - 1/3\) and \(w(C, \phi) = \sum_{p=1}^{m} w(C_p, \phi)\).

Lemma 3.2. If \(w(C, \phi) \geq \kappa\) then \((V, C)\) is a Yes-instance of \text{BATLB}.

Proof. By linearity of expectation, \(w(C, \phi) \geq \kappa\) implies \(\mathbb{E}\left(\sum_{p=1}^{m} v_p(\alpha)\right) \geq m/3 + \kappa\). Thus, if \(w(C, \phi) \geq \kappa\) then there is a \(\phi\)-compatible bijection \(\alpha\) that satisfies at least \(m/3 + \kappa\) constraints. \(\square\)

Let \(X = w(C, \phi)\) and \(X_p = w(C_p, \phi), p = 1, \ldots, m\). Observe that if \(\phi\) is a random function from \(V\) to \(\{0, 1, 2, 3\}\) then \(X, X_1, \ldots, X_m\) are random variables. Recall that \(X = \sum_{p=1}^{m} X_p\).

Lemma 3.3. We have \(\mathbb{E}[X] = 0\).

Proof. Let \(C_p = (v_i, \{v_j, v_k\}) \in C\). Let us first find the distribution of \(X_p\). It is easy to check that the probability that \(\phi(v_i) = \phi(v_j) = \phi(v_k)\) equals 1/16 and \(X_p = 0\) in such a case. The probability that \(\phi(v_i) \neq \phi(v_j) = \phi(v_k)\) equals 3/16 and \(X_p = -1/3\) in such a case. The probability that \(\phi(v_i)\) equals one of the non-equal \(\phi(v_j), \phi(v_k)\) is equal to 6/16 and \(X_p = 1/6\) in such a case. Now suppose that \(\phi(v_j), \phi(v_j)\) and \(\phi(v_k)\) are all distinct. The probability that \(\phi(v_i)\) is between \(\phi(v_j)\) and \(\phi(v_k)\) is 2/16 and \(X_p = 2/3\) in such a case. Finally, the probability that \(\phi(v_i)\) is not between \(\phi(v_j)\) and \(\phi(v_k)\) is 4/16 and \(X_p = -1/3\) in such a case. Now we can give the distribution of \(X_p\) in the following table.

Using this distribution, it is easy to see that \(\mathbb{E}[X] = 0\) and, thus, \(\mathbb{E}[X] = \sum_{p=1}^{m} \mathbb{E}[X_p] = 0\). \(\square\)
For an irreducible instance Lemma 3.5. The above polynomial is of degree 6. It remains to recall \(X\) will compute \(E[X]\) of 0. Indeed, (1) and \(E[X] = 11/96 = 88/768\). It remains to show that

\[
\sum_{1 \leq l \neq l' \leq m} E[X_l X_{l'}] \geq -\frac{77}{768} m.
\] (1)

Indeed, (1) and \(E[X^2] = 88/768\) imply that

\[
E[X^2] = \sum_{l=1}^{m} E[X_l^2] + \sum_{1 \leq l \neq l' \leq m} E[X_l X_{l'}] \geq \frac{88}{768} m - \frac{77}{768} m = \frac{11}{768} m,
\]
In the reminder of this proof we show that (1) holds. Let $C_{i}, C_{\ell'}$ be a pair of distinct constraints of $C$. To evaluate $E[X_{i}X_{\ell'}]$, we consider several cases. A simple case is when the sets $vars(C_{i})$ and $vars(C_{\ell'})$ are disjoint: then $X_{i}$ and $X_{\ell'}$ are independent random variables and, thus, $E[X_{i}X_{\ell'}] = E[X_{i}]E[X_{\ell'}] = 0$. Let $U = \{(l,l') \mid C_{i}, C_{\ell'} \in C, l \neq l'\}$ be the set of all ordered index pairs corresponding to distinct constraints in $C$. We will classify subcases of this case by considering some subsets of $U$. Let

$$S_{1}(u) = \{(l,l') \in U \mid C_{i} = (u,\{a,b\}), C_{\ell'} = (u,\{c,d\}), a, b, c, d \in V\},$$

$$S_{2}(u) = \{(l,l') \in U \mid C_{i} = (u,\{a,b\}), C_{\ell'} = (c,\{u,d\}), a, b, c, d \in V\},$$

$$S_{3}(u) = \{(l,l'), (l',l) \in U \mid C_{i} = (u,\{a,b\}), C_{\ell'} = (c,\{u,d\}), a, b, c, d \in V\},$$

$$S_{4}(u,v) = \{(l,l') \in U \mid C_{i} = (u,\{v,a\}), C_{\ell'} = (u,\{v,b\}), a, b, v, d \in V\},$$

$$S_{5}(u,v) = \{(l,l') \in U \mid C_{i} = (u,\{v,a\}), C_{\ell'} = (c,\{u,v\}), a, b, v, d \in V\},$$

$$S_{6}(u,v) = \{(l,l'), (l',l) \in U \mid C_{i} = (u,\{v,a\}), C_{\ell'} = (c,\{u,v\}), a, b, v, d \in V\},$$

$$S_{7}(u,v, w) = \{(l,l') \in U \mid \text{vars}(C_{i}) = \text{vars}(C_{\ell'}) = \{u, v, w\}\}.$$

Let $u, v \in V$ be a pair of distinct variables. Observe that $S_{4}(u,v) = (S_{1}(u) \cap S_{2}(v)) \cup (S_{1}(v) \cap S_{2}(u))$, $S_{5}(u,v) = S_{2}(u) \cap S_{4}(v)$, $S_{6}(u,v) = (S_{1}(u) \cap S_{2}(v)) \cup (S_{3}(v) \cap S_{2}(u))$ and $S_{7}(u,v) = S_{3}(u) \cap S_{3}(v)$. Let $u, v, w \in V$ be a triple of distinct variables. Observe that

$$S_{6}(u,v,w) = (S_{3}(u) \cap S_{3}(v) \cap S_{2}(w)) \cup (S_{3}(v) \cap S_{3}(u) \cap S_{2}(w)) \cup (S_{3}(w) \cap S_{3}(u) \cap S_{2}(v)).$$

(2)

For a variable $u \in V$, let $b(u) = |\{l : C_{i} = (u,\{a,b\}), a, b \in V\}|$ and $e(u) = |\{l : C_{i} = (u,\{a,b\}), a, b \in V\}|$. Observe that $|S_{1}(u)| = b(u)(b(u) - 1)$, $|S_{2}(u)| = e(u)(e(u) - 1)$ and $|S_{3}(u)| = 2b(u)e(u)$. For a pair $u, v \in V$, let $c_{uv}^{n} = |\{l : C_{i} = (u,\{v,a\}), a \in V\}|$ and $c_{uv}^{n} = |\{l : C_{i} = (u,\{v,a\}), a \in V\}|$. Observe that $|S_{4}(u,v)| = c_{uv}^{n}b(u)(b(u) - 1) + c_{uv}^{n}e(u)(e(u) - 1)$, $|S_{5}(u,v)| = c_{uv}^{n}c_{uv}^{n} - 1), |S_{6}(u,v)| = 2(c_{uv}^{n} + c_{uv}^{n}) \cdot c_{uv}^{n}$ and $|S_{7}(u,v)| = 2c_{uv}^{n}c_{uv}^{n}$. Let $u, v, w \in V$ be a triple of distinct variables. Since $C$ is irreducible, the number of ordered pairs $(C_{i}, C_{\ell'})$ for which $\text{vars}(C_{i}) = \text{vars}(C_{\ell'}) = \{u, v, w\}$ is at most 2, i.e., $|S_{6}(u,v,w)| \leq 2$.

We list the sets $S_{i}(\cdot)$, their union/intersection forms (for $i = 4, 5, 6, 7$) and their sizes in Table 2. If $(l,l')$ belongs to some $S_{i}$ but to no $S_{j}$ for $j > i$, then Table 2 also contains the value $768 \cdot E[X_{i}X_{\ell'}]$, in the row corresponding to $S_{i}$. These values cannot be easily calculated analytically as there are many cases to consider and we have calculated them using computer. We will briefly describe how our program computes $E[X_{i}X_{\ell'}]$ using as an example the case $(l,l') \in S_{1}(u), i.e., C_{i} = (u,\{a,b\}), C_{\ell'} = (u,\{c,d\})$. For each $(q_{1}, q_{2}, q_{3}, q_{4}, q_{5}) \in \{0, 1, 2, 3\}^{5}$ the probability of $(u, a, b, c, d) = (q_{1}, q_{2}, q_{3}, q_{4}, q_{5})$ is $4^{-5}$ and the corresponding value of $X_{i}X_{\ell'}$ can be found in Table 2.
We are now ready to compute a lower bound on the term $\sum_{1 \leq l \neq l' \leq m} E[XX_t]$. Define the values $w'_i$ for $i = 1, 2, \ldots, 8$ as it is done in Table 2. We will now show that the following holds (note that the sets we sum over have to contain distinct elements).

$$\sum_{1 \leq l \neq l' \leq m} E[XX_t] = \sum_{u \in V} \sum_{i=1}^{3} |S_i(u)| w'_i + \sum_{\{u, v\} \subseteq V} \sum_{i=4}^{7} |S_i(u, v)| w'_i + \sum_{\{u, v, w\} \subseteq V} |S_8(u, v)| w'_8$$

In order to show the above we consider the possible cases for $(l, l') \in U$.

**Case 1**: $|\text{vars}(C_l) \cap \text{vars}(C_{l'})| = 0$. In this case $E[XX_t] = 0$ and the corresponding $(l, l')$ does not belong to any $S_i$ and therefore contributes zero to the right-hand side above.

**Case 2**: $|\text{vars}(C_l) \cap \text{vars}(C_{l'})| = 1$. Each pair $(l, l') \in S_1(u)$ contributes $\frac{12}{768}$ to both sides of the above equation, as in this case $(l, l')$ does not belong to any $S_j$ with $j > 1$. Analogously if $(l, l') \in S_2(u)$ then it contributes $\frac{3}{768}$ to both sides of the above equation. Furthermore if $(l, l') \in S_3(u)$ then it contributes $-\frac{6}{768}$.

**Case 3**: $|\text{vars}(C_l) \cap \text{vars}(C_{l'})| = 2$. Consider a pair $(l, l') \in S_4(u, v)$ and assume, without loss of generality, that $(l, l') \in S_1(u) \cap S_2(v)$. Note that $(l, l')$ contributes $\frac{31}{768}$ to the left-hand side of the equation and it contributes $w'_5 + w'_6 + w'_7 = \frac{31}{768}$ to the right-hand side as $(l, l') \in S_1(u) \cap S_2(v) \cap S_4(u, v))$. Analogously if $(l, l') \in S_5(u, v)$ we get a contribution of $w'_5 = \frac{36}{768} = w'_3 + w'_2 + w'_1$ to both sides of the equation. If $(l, l') \in S_6(u, v)$ we get a contribution of $w'_6 = -\frac{15}{768} = w'_3 + w'_2 + w'_1$ to both sides of the equation. If $(l, l') \in S_7(u, v)$ we get a contribution of $w'_7 = -\frac{11}{768} = w'_3 + w'_2 + w'_1$ to both sides of the equation.

**Case 4**: $|\text{vars}(C_l) \cap \text{vars}(C_{l'})| = 3$. Assume, without loss of generality, that $(l, l') \in S_3(u) \cap S_3(v) \cap S_2(w)$ and note that $(l, l') \in S_2(u, v) \cap S_6(u, w) \cap S_8(v, w)$.

### Table 2: Data for sets $S_i(u)$, $i = 1, 2, \ldots, 8$.

| Set | Union/intersection Form | $|\text{Set}|$ | $768E[XX_t]$ | $768w'_i$ |
|-----|-------------------------|--------------|--------------|-----------|
| $S_1(u)$ | $-$ | $b(u)(b(u) - 1)$ | 12 = $w_1$ | 12 |
| $S_2(u)$ | $-$ | $c(u)(c(u) - 1)$ | 3 = $w_2$ | 3 |
| $S_3(u)$ | $-$ | $b(u)c(u) + e(u)b(u)$ | $-6 = w_3$ | $-6$ |
| $S_4(u, v)$ | $(S_1(u) \cap S_2(v)) \cup (S_1(v) \cap S_2(u))$ | $c_u(e_u - 1) + c_v(e_v - 1)$ | 24 = $w_4$ | 9 |
| $S_5(u, v)$ | $S_2(u) \cap S_2(v)$ | $c_{uv}(e_{uv} - 1)$ | 36 = $w_5$ | 30 |
| $S_6(u, v)$ | $(S_3(u) \cap S_2(v)) \cup (S_3(v) \cap S_2(u))$ | $2(c_v + c_u) \cdot c_{uv}$ | $-18 = w_6$ | $-15$ |
| $S_7(u, v)$ | $S_3(u) \cap S_3(v)$ | $2c_v c_u^2$ | $-6 = w_7$ | 6 |
| $S_8(u, v, w)$ | see (2) | $\leq 2$ | $-44 = w_8$ | $-11$ |
Therefore we get a contribution of \( w_8 = \frac{-44}{768} = w'_5 + w'_4 + w'_2 + w'_6 + w'_8 \) to both sides of the equation.

Therefore the above equation holds, which implies the following:

\[
\sum_{1 \leq i \neq i' \leq m} \mathbb{E}[X_i X_{i'}] = \sum_{u \in V} (|S_1(u)|w'_i + |S_2(u)|w'_i + |S_3(u)|w'_i) \\
+ \sum_{(u,v) \subseteq V} \sum_{i=1}^{7} |S_i(u,v)|w'_i + \sum_{(u,v,w) \subseteq V} |S_8(u,v,w)|w'_8
\]

\[
= \frac{1}{2 \cdot 768} \sum_{u \in V} \left( 6(2b(u) - e(u))^2 - 24b(u) - 6e(u) \right) \\
+ \frac{1}{2 \cdot 768} \sum_{(u,v) \subseteq V} \left( 15(e'_v + c'_v - 2c'_{uv})^2 + 12 \left( \frac{c'_v - c'_u}{2} \right)^2 - 18(c'_v + c'_u) - 60c'_{uv} \right) \\
+ \sum_{(u,v,w) \subseteq V} |S_8(u,v,w)|w'_8
\]

To complete the proof of the lemma it remains to translate this sum into a function on the number of constraints. In that respect, notice that \( \sum_{u \in V} b(u) = m \) and \( \sum_{u \in V} e(u) = 2m \). Further, each clause \( (u, \{v, w\}) \) contributes exactly one unit to each of \( c'_v \) and \( c'_w \), as well as exactly one unit to \( c'_{uv} \). Hence \( \sum_{(u,v) \subseteq V} (c'_v + c'_w) = 2m \) and \( \sum_{(u,v,w) \subseteq V} c'_{uv} = m \). Since \( \mathcal{C} \) is irreducible, the number of ordered pairs \( (C, C') \) for which \( \text{vars}(C_i) = \text{vars}(C_i') \) is at most \( m/2 \) and, thus,

\[
\sum_{(u,v,w) \subseteq V} |S_8(u,v,w)|w'_8 \leq m \cdot w'_8.
\]

Together these bounds imply that

\[
\sum_{1 \leq i \neq i' \leq m} \mathbb{E}[X_i X_{i'}] \geq -\frac{36}{2 \cdot 768} m - \frac{96}{2 \cdot 768} m - \frac{11}{768} m - \frac{77}{768} m
\]

and \( \Box \) holds.

We are now ready to prove the main result.

**Theorem 3.6.** BATLB has a kernel of size \( O(\kappa^2) \).

**Proof.** Let \((V, \mathcal{C})\) be an instance of BATLB. By Lemma 3.1 in time \( O(m^3) \) we can obtain an irreducible instance \((V', \mathcal{C}')\) such that \((V, \mathcal{C})\) is a Yes-instance if and only if \((V', \mathcal{C}')\) is a Yes-instance. Let \( m' = |\mathcal{C}'| \) and let \( X \) be the random variable defined above. Then \( X \) is expressible as a polynomial of degree 6 by Lemma 3.4, hence it follows from Lemma 2.2 that \( \mathbb{E}[X^4] \leq 2^{2m} \mathbb{E}[X^2]^2 \). Consequently, \( X \) satisfies the conditions of Lemma 2.1 from which we conclude in combination with Lemma 3.5 that \( \mathbb{P} \left( X > \frac{1}{4 \pi^2} \sqrt{\frac{11}{768}} m' \right) > 0 \). By Lemma 3.2 if \( \frac{11}{768} m' \geq \kappa \) then \((V', \mathcal{C}')\) is a Yes-instance for BATLB. Otherwise, we have \( m' = O(\kappa^2) \). This concludes the proof of the theorem. \( \Box \)
We complete this section by answering the following natural question: why have we considered functions \( \phi : V \to \{0, 1, 2, 3\} \) rather than functions \( \phi : V \to \{0, 1\} \)? The latter would involve less computations and give a smaller degree of the polynomial representing \( X \). The reason is that our proof of Lemma 3.5 would not work for functions \( \phi : V \to \{0, 1\} \) (we would only be able to prove that \( E[X^2] \geq \sum_{\{u, v\} \subset V} [c_u^v + c_v^u - 2c_{uv}]^2 \), which is not enough).

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