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Computing groups of Hecke characters

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Abstract

We describe algorithms to represent and compute groups of Hecke characters. We make use of an idèlic point of view and obtain the whole family of such characters, including transcendental ones. We also show how to isolate the algebraic characters, which are of particular interest in number theory. This work has been implemented in Pari/GP, and we illustrate our work with a variety of explicit examples using our implementation.

1 Introduction

Hecke characters are, from the modern point of view, continuous characters of idèle class groups, in other words automorphic forms for \( \text{GL}_1 \). They were introduced by Hecke [13] who proved the functional equation of their \( L \)-function, and are the starting point of many developments that blossom in modern number theory: automorphic \( L \)-functions via Tate’s thesis [39], \( \ell \)-adic Galois representations via Weil’s notion of algebraic characters [43], Shimura varieties via CM theory [38], and the Langlands programme via class field theory and the global Weil group [44]. Despite their fundamental role, Hecke characters have not received a full algorithmic treatment, perhaps due to the fact that they are considered well-understood compared to automorphic forms on higher rank groups. The existing literature only describes how to compute with finite order characters, since they are characters of ray class groups [7], and algebraic Hecke characters [42]. As part of a collective effort to enumerate and compute \( L \)-functions, automorphic representations and Galois representations, we believe that the \( \text{GL}_1 \) case also deserves close scrutiny, and this is the goal of the present paper.

We describe algorithms to compute, given a number field \( F \) and a modulus \( m \) over \( F \), a basis of the group of Hecke quasi-characters of modulus \( m \).
(Algorithm 18) and its subgroup of algebraic characters (Algorithm 30), in a form suitable for evaluation at arbitrary ideals and decomposition into local characters (Algorithm 17). In particular, we describe a polynomial time algorithm to compute the maximal CM subfield of $F$ (Algorithm 28). It is sometimes believed that the adèlic point of view is not suitable for computational purposes; we claim the contrary, and adopt an adèlic setting throughout the paper. Our implementation in Pari/GP is available from version 2.15 of the software. We provide examples that illustrate the use of our algorithms and showcase some interesting features of Hecke characters: a presentation of the software interface, small degree examples, illustrations of automorphic induction from quadratic fields, examples of CM abelian varieties with emphasis on the rigorous identification of the corresponding Hecke character, illustration of the density of the gamma shifts of Hecke $L$-functions in the conjectured space of possible ones (Proposition 44), examples of provably partially algebraic Hecke characters (Proposition 46) and of twists of $L$-functions by Hecke characters.

The only previous work on computation of infinite order Hecke characters is that of Watkins, so we give a short comparison: in Watkins’s paper, only algebraic characters were considered, and only over a CM field, whereas we treat arbitrary Hecke characters over arbitrary number fields; the values of characters were represented exactly by algebraic numbers, whereas we represent values by approximations since this is forced in the transcendental case; the emphasis was on individual Hecke characters, which the user had to construct by hand, whereas our emphasis is on groups of Hecke characters, which we construct for the user, simply from the modulus.

Our implementation makes it possible to tabulate Hecke characters and their $L$-functions systematically by increasing analytic conductor; we think that this is a valuable project but we leave it for future work.

The paper is organized as follows. In Section 2 we recall the definitions and basic properties of Hecke characters and their $L$-functions. In Section 3 we describe our algorithms to compute groups of Hecke characters and evaluate them. In Section 4 we present our algorithms to compute the maximal CM subfield and groups of algebraic Hecke characters. Finally, Section 5 contains a variety of examples.

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2 Hecke characters

We recall the definition of Hecke characters in the adèlic setting. This material is standard and can be found in \[18\] chap. XIV or \[34\].

Let \(F\) be a number field of degree \([F : \mathbb{Q}] = n\) and discriminant \(\Delta_F\). When \(K/F\) is a finite extension, we denote by \(N_{K/F}\) the norm from \(K\) to \(F\); we also denote \(N = N_{F/\mathbb{Q}}\) when \(F\) is clear from the context. For every prime ideal \(\mathfrak{p}\) of \(F\), we consider the completion \(F_{\mathfrak{p}}\) and its ring of integers \(\mathcal{O}_{F_{\mathfrak{p}}}\). We choose a uniformizer \(\pi_{\mathfrak{p}} \in \mathcal{O}_{F_{\mathfrak{p}}}\) and denote by \(v_{\mathfrak{p}}: F_{\mathfrak{p}} \times \mathbb{F}_{\mathfrak{p}} \to \mathbb{Z}\) the \(\mathfrak{p}\)-adic valuation. We will always use \(\sigma\) to denote an archimedean place of \(F\) and the corresponding real or complex embedding. For every place \(v\), let \(n_v = [F_v : \mathbb{Q}_v]\), and let \(\cdot | v\) be the normalized absolute value, i.e. \(n_\sigma = 1\) and \(|\cdot|_\sigma = |\cdot|\) for a real embedding \(\sigma\), \(n_\sigma = 2\) and \(|\cdot|_\sigma = |\cdot|^2\) for a complex embedding \(\sigma\), and \(|\pi_{\mathfrak{p}}|_{\mathfrak{p}} = N_p^{-1}\) for a prime ideal \(\mathfrak{p}\). We denote by \(A^\times_F = \prod' F_v^\times\) the group of idèles of \(F\). We write \(F_{\mathbb{R}} = F \otimes_{\mathbb{Q}} \mathbb{R} \cong \prod_\sigma F_\sigma \cong \mathbb{R}^{r_1} \times \mathbb{C}^{r_2}\), where \(r_1\) (resp. \(r_2\)) is the number of real embeddings (resp. pairs of non-real complex embeddings) of \(F\).

Let \(U\) denote the group of complex numbers of absolute value 1. For \(G\) a topological group, \(G^\circ\) will denote the connected component of 1 in \(G\).

2.1 Pontryagin duality

We recall some definitions and properties of locally compact abelian groups that will be used later. See \[28, 29\] for general reference.

Let \(G\) be a locally compact abelian group. A quasi-character of \(G\) is a continuous morphism
\[
\chi: G \to \mathbb{C}^\times.
\]
A character of \(G\) is a continuous morphism
\[
\chi: G \to U.
\]

The group of characters of \(G\), which we denote by \(\hat{G}\), is the Pontryagin dual \(\text{Hom}_{\text{cont}}(G, U)\) of \(G\), and is a locally compact abelian group. The canonical map
\[
G \to \hat{G}
\]
given by \(g \mapsto (\chi \mapsto \chi(g))\) is an isomorphism. Let \(H \subset G\) be a subgroup. Let
\[
H^\perp = \{\chi \in \hat{G} \mid \chi(h) = 1 \text{ for all } h \in H\}
\]
be the Pontryagin orthogonal of \(H\) in \(\hat{G}\). Then \(H^\perp\) is a closed subgroup of \(\hat{G}\), and \((H^\perp)^\perp\) is the closure of \(H\), where the second orthogonal is taken in \(G\). If \(H\) is a closed subgroup of \(G\), then we have canonical isomorphisms
\[
\hat{G}/H \cong H^\perp \text{ and } \hat{G}/(H^\perp)^\perp \cong \hat{H}.
\]
The group $G$ is compact if and only if $\hat{G}$ is discrete.

Pontryagin duality is an exact contravariant functor on the category of locally compact abelian groups.

Let $(x, y) \mapsto x \cdot y$ denote a nondegenerate $\mathbb{R}$-bilinear form on a finite dimensional $\mathbb{R}$-vector space $V$. The pairing $V \times V \to \mathbb{U}$ defined by $(x, y) \mapsto \exp(2\pi i x \cdot y)$ induces an isomorphism $V \cong \hat{V}$. We will use this isomorphism to identify characters on $V$ with elements of $V$.

Let $\Lambda$ be a full rank lattice in $V$. The pairing above identifies the dual lattice $\Lambda^\perp = \text{Hom}(\Lambda, \mathbb{Z})$ with the subgroup $\Lambda^\perp = \{x \in V \mid x \cdot y \in \mathbb{Z} \text{ for all } y \in \Lambda\}$, which is canonically isomorphic to $\hat{V}/\Lambda$ by the above, and we have $\hat{\Lambda} \cong V/\Lambda^\perp$. In particular for $V = \mathbb{R}$ and $\Lambda = \mathbb{Z}$ we consider the standard bilinear form and we have $\hat{\mathbb{R}}/\mathbb{Z} = \mathbb{Z}^\perp = \mathbb{Z}$ and $\hat{\mathbb{Z}} = \mathbb{R}/\mathbb{Z}$.

The dual $V = \hat{\mathbb{Q}}$ of the group of rationals equipped with the discrete topology, is the compact topological group $\varprojlim \mathbb{R}/n\mathbb{Z}$, called the solenoid.

### 2.2 General Hecke characters

A Hecke quasi-character is a quasi-character of $C_F = \mathbb{A}_F^\times/F^\times$, and a Hecke character is a character of $C_F$.

The *norm* is the Hecke quasi-character

$$\| \cdot \| : C_F \to \mathbb{C}^\times$$

defined by

$$x = (x_v) \mapsto \|x\| = \prod_v |x_v|_v.$$

This is a well-defined Hecke quasi-character by the product formula.

Every Hecke quasi-character $\chi$ is of the form $\chi = \chi_0 \| \cdot \|^s$ for a unique Hecke character $\chi_0$ and a unique $s \in \mathbb{R}$. We refer to $\chi_0$ as the *unitary component* of $\chi$. In the algebraic setting, the value $w = -2s$ is the weight of $\chi$.

We also define $C_F^1 = \ker(\| \cdot \| : C_F \to \mathbb{R}_{>0})$ to be the kernel of the norm, which is a compact group. We have a canonical embedding

$$\mathbb{R}_{>0} \to C_F,$$

by sending $t \in \mathbb{R}_{>0} \mapsto ((t^{1/n})_\sigma, 1, \ldots) \in \mathbb{A}_F^\times$ where $t \mapsto (t^{1/n})_\sigma$ denotes the diagonal embedding $\mathbb{R}_{>0} \to \prod_\sigma F_\sigma^\times$, and a canonical decomposition

$$C_F \cong C_F^1 \times \mathbb{R}_{>0}.$$

As a consequence, it suffices to compute the characters of $C_F^1$ to deduce the full groups of Hecke characters and Hecke quasi-characters

$$\text{Hom}_{\text{cont}}(C_F, \mathbb{C}^\times) = \hat{C}_F \| \cdot \| \cong \hat{C}_F^1 \| \cdot \|^\mathbb{C}. \quad (1)$$
Every quasi-character $\chi$ of $A_F^\times$ (and in particular every Hecke quasi-character) admits a factorization $\chi = \prod_v \chi_v$, where $\chi_v$ is a quasi-character of $F_v^\times$. We therefore describe quasi-characters of local fields.

### 2.3 Local characters

- Every quasi-character $\chi$ of $C^\times$ is of the form
  \[
  \chi(z) = \left(\frac{z}{|z|^k}\right)^k |z|^{2s}
  \]
  for a unique pair $(k, s) \in \mathbb{Z} \times \mathbb{C}$. The quasi-character $\chi$ is a character if and only if $\Re(s) = 0$, i.e. $s = i\varphi$ for some $\varphi \in \mathbb{R}$.

- Every quasi-character $\chi$ of $R^\times$ is of the form
  \[
  \chi(x) = \text{sgn}(x)^{k|x|^s}
  \]
  for a unique pair $(k, s) \in \{0, 1\} \times \mathbb{C}$. We say that $\chi$ is unramified if $k = 0$. The quasi-character $\chi$ is a character if and only if $\Re(s) = 0$, i.e. $s = i\varphi$ for some $\varphi \in \mathbb{R}$.

- Let $p$ be a prime ideal of $\mathbb{Z}_F$. Every quasi-character $\chi$ of $F_p^\times$ is of the form
  \[
  \chi(x) = \chi_0(x\pi_p^{-r_p(x)} \mod p^m)\chi(p)^{r_p(x)}
  \]
  for a unique $m \geq 0$ and a unique primitive character $\chi_0$ of $(\mathbb{Z}_p/p^m)^\times$, and where we write $\chi(p) = \chi(p_\sigma)$ of $C^\times$. Note that in general $\chi(p)$ depends on the choice of uniformizer $\pi_p$, but $\chi(p)$ is well defined up to the roots of unity of the same order as $\chi_0$. We call $p^m$ the conductor of $\chi$ and $m$ its conductor exponent. If $m = 0$ we call $\chi$ unramified; in this case, $\chi(p)$ does not depend on the choice of uniformizer, and the quasi-character $\chi$ only depends on $\chi(p)$. Regardless of $m$, the quasi-character $\chi$ is a character if and only if $\chi(p) \in U$.

Whenever we write a global idèle character $\chi$ as a product of local characters $\chi_v$, we write its local parameters $k_\sigma, \varphi_\sigma$, and $m_\sigma$, and we let $f_\chi = \prod_p p^{m_\sigma}$ be the conductor of $\chi$. Note that for a complex place, the pair $(k_\sigma, \varphi_\sigma)$ depends on the choice of a complex embedding among the two conjugate ones, or equivalently on the choice of an isomorphism between the completion of $F$ and $\mathbb{C}$; we have $\varphi_\overline{\sigma} = \varphi_\sigma$ and $k_\overline{\sigma} = -k_\sigma$.

### 2.4 $L$-function

Let $\chi$ be a Hecke character such that $\sum_\sigma n_\sigma \varphi_\sigma = 0$, i.e. that is trivial on the embedded $\mathbb{R}_{>0}$ in $\mathbb{A}$. Let $N_\chi = |\Delta_F| \cdot N(f_\chi)$. Let

\[
L(\chi, s) = \prod_{\mathfrak{p} \nmid f_\chi} (1 - \chi(\mathfrak{p}) N(\mathfrak{p})^{-s})^{-1}
\]
and
\[
\gamma(\chi, s) = \prod_{\sigma \text{ real}} \Gamma_R(s + i\varphi_{\sigma} + k_{\sigma}) \cdot \prod_{\sigma \text{ complex}} \Gamma_C(s + i\varphi_{\sigma} + |k_{\sigma}|/2).
\]
where \( \Gamma_R(s) = \pi^{-s/2} \Gamma(s^2) \) and \( \Gamma_C(s) = 2(2\pi)^{-s} \Gamma(s) \) Then
\[
\Lambda(\chi, s) = N_s^{s/2} \chi^{\gamma(\chi, s)} L(\chi, s)
\]
satisfies the functional equation
\[
\Lambda(\chi, 1 - s) = W(\chi) \Lambda(\bar{\chi}, s)
\]
for some complex number \( W(\chi) \) of absolute value 1.

We have the formula
\[
W(\chi) = \prod_v W(\chi_v),
\]
where
\[
W(\chi_v) = \begin{cases} 
4^{i\varphi_v} v_{\sigma}^{|k_{\sigma}|} & \text{if } v = \sigma \text{ is complex}, \\
i^{k_{\sigma}} & \text{if } v = \sigma \text{ is real}, \\
\chi(p)^{d_{p\sigma}} \chi(p) \eta(\chi_p) N(p)^{-m_p/2} & \text{if } v = p | f_{\chi}, \text{ and} \\
\chi(p)^{d_{p\sigma}} & \text{if } v = p \not{|} f_{\chi}.
\end{cases}
\]
where \( d_{p\sigma} = v_p(\mathfrak{D} f_{\chi}) \) and \( \mathfrak{D} \) is the different of \( F \) (so that the product is finite), and
\[
\tau(\chi_p) = \sum_{\epsilon \in \mathcal{O}_F^\times} \chi(\epsilon) \exp(2i\pi \lambda \circ \text{Tr}_{F_p/Q}(\epsilon/fp^\mu))
\]
where \( \lambda : Q_p \to Q_p/Z_p \to Q/Z \).

### 2.5 Algebraic Hecke characters

**Warning:** an algebraic Hecke character is usually not a Hecke character, it is only a quasi-character.

Let \( \chi \) be a Hecke quasi-character. It is called algebraic if for every archimedean place \( \sigma \) of \( F \), there exists integers \( p_{\sigma}, q_{\sigma} \in \mathbb{Z} \) such that for all \( z \in (F_{\sigma}^\times)^o \) we have\(^1\)
\[
\chi_{\sigma}(z) = z^{-p_{\sigma}}(z)^{-q_{\sigma}}.
\]

Note: if \( \sigma \) is complex, then \( p_{\sigma} \) and \( q_{\sigma} \) are uniquely determined; if \( \sigma \) is real then only their sum is well-defined. We say that \( \chi \) is of type \( (p_{\sigma}, q_{\sigma})_{\sigma} \).

\(^1\)The choice of sign in the exponents is such that the values of \( \chi \) at integral ideals are algebraic integers if and only if all \( p_{\sigma} \) and \( q_{\sigma} \) are nonnegative.
Example 1. The norm $\| \cdot \|$ is an algebraic character, of type $(p_\sigma, q_\sigma) = (-1, -1)$ if $\sigma$ is complex. We have $\|p\| = N(p)^{-1}$ for every prime ideal $p$.

Definition 2. We call a Hecke character almost-algebraic if $\varphi_\sigma = 0$ for all $\sigma$. We denote by $(\hat{C}_F)^{\text{aa}}$ the subgroup of almost-algebraic characters.

Remark 3. Algebraic characters correspond to type $A_0$ and almost-algebraic to type $A$ with trivial norm component in Weil’s terminology [43]. By a theorem of Waldschmidt [41], these definitions coincide with the fact that a quasi-character has type $A$ if and only if its values are algebraic, and type $A_0$ if and only if there exists a finite extension of $\mathbb{Q}$ containing all of its values.

2.5.1 Parameters at infinity of algebraic Hecke characters

It is known that if $F$ has a real embedding, then every algebraic Hecke character is an integral power of the norm times a Hecke character of finite order (see [43]). So from now on we assume that $F$ is totally complex. We recall the following well-known lemma.

Lemma 4. Let $\chi_0$ be a Hecke character and let $(k_\sigma, \varphi_\sigma)$ denote its local parameters at infinite places. The character $\chi_0$ is the unitary component of an algebraic Hecke character if and only if $\chi_0$ is almost algebraic and all $k_\sigma$ have the same parity.

More precisely, let $\chi = \chi_0 \| \cdot \|^{-w/2}$ be a Hecke quasi-character with $w \in \mathbb{R}$. If $\chi$ is algebraic of type $(p_\sigma, q_\sigma)$, then

- $w \in \mathbb{Z}$;
- $p_\sigma + q_\sigma = w$ for all $\sigma$;
- $k_\sigma = q_\sigma - p_\sigma$ for all $\sigma$;
- $\varphi_\sigma = 0$ for all $\sigma$.

Conversely, if $\chi_0$ is almost-algebraic and all $k_\sigma$ have the same parity, let $w \in \mathbb{Z}$ have the same parity as the $k_\sigma$; then $\chi = \chi_0 \| \cdot \|^{-w/2}$ is algebraic.

Proof. Let $\chi = \chi_0 \| \cdot \|^{-w/2}$ be a Hecke quasi-character with $w \in \mathbb{R}$, so that for all $z \in \mathbb{C}^\times$ we have

$$\chi_\sigma(z) = \left( \frac{z}{|z|} \right)^{k_\sigma} |z|^{2i\varphi_\sigma - w}.$$

Let $p, q \in \mathbb{Z}$. For all $z \in \mathbb{C}^\times$ we have

$$z^{-p}(\bar{z})^{-q} = \left( \frac{z}{|z|} \right)^{q-p} |z|^{-p-q}.$$

By uniqueness of parameters of quasi-characters of $\mathbb{C}^\times$, the quasi-character $\chi$ is algebraic of type $(p_\sigma, q_\sigma)$ if and only if for all $\sigma$ we have $k_\sigma = q_\sigma - p_\sigma$,
\[ \varphi_\sigma = 0 \text{ and } w = p_\sigma + q_\sigma. \] In this case, \( \chi_0 \) is almost-algebraic and for all \( \sigma \) we have \( k_\sigma \equiv q_\sigma - p_\sigma \equiv p_\sigma + q_\sigma \equiv w \mod 2 \), so that all \( k_\sigma \) have the same parity. This also validates the construction of an algebraic \( \chi \) from an almost-algebraic \( \chi_0 \) satisfying the parity condition.

Thus the group of unitary components \( \chi_0 \) of algebraic Hecke characters \( \chi = \chi_0 \cdot \| - w/2 \) is a finite index subgroup of the group of almost-algebraic Hecke characters.

### 2.5.2 \( L \)-function of an algebraic Hecke character

Let \( \chi = \chi_0 \cdot \| - w/2 \) be an algebraic Hecke character as above. Let \( f_\chi = f_{\chi_0} \) be its conductor and \( N_\chi = N_{\chi_0} \). Let

\[
L(\chi, s) = \prod_{p \nmid f_\chi} \left((1 - \chi(p) N(p)^{-s})^{-1} = L(\chi_0, s - w/2),
\right.
\]

and

\[
\gamma(\chi, s) = \prod_\sigma \Gamma_C(s - \min(p_\sigma, q_\sigma)) = \gamma(\chi_0, s - w/2).
\]

Then

\[
\Lambda(\chi, s) = N^{s/2}_\chi \gamma(\chi, s) L(\chi, s)
\]

satisfies the functional equation

\[
\Lambda(\chi, w + 1 - s) = W(\chi) \Lambda(\overline{\chi}, s)
\]

for some complex number \( W(\chi) = W(\chi_0) \) of absolute value 1.

### 3 Computing the group of Hecke characters

#### 3.1 Filtration by modulus

We have a non-canonical isomorphism

\[
\widehat{C}_F \cong \mathbb{Q}^{r_1 + r_2 -1} \times \mathbb{Z}_r^2 \times \mathbb{R},
\]

where \( T \) is an infinite torsion abelian group. Indeed, we have the classical decomposition [44]

\[
1 \rightarrow C_F^0 \rightarrow C_F \rightarrow \pi_0(C_F) \rightarrow 1, \text{ where } C_F^0 \cong \mathbb{V}^{r_1 + r_2 -1} \times (\mathbb{R}/\mathbb{Z})^{r_2} \times \mathbb{R},
\]

where \( \mathbb{V} = \widehat{\mathbb{Q}} \) is the solenoid, and \( \pi_0(C_F) \) is profinite; by Pontryagin duality, we get

\[
0 \rightarrow T \rightarrow \widehat{C}_F \rightarrow \mathbb{Q}^{r_1 + r_2 -1} \times \mathbb{Z}_r^2 \times \mathbb{R} \rightarrow 0,
\]

and this exact sequence splits. Since we cannot give a finite description of the whole group \( T \), we will filter \( \widehat{C}_F \) according to moduli.
Let \( m = m_f m_\infty \) be a modulus, meaning that \( m_f \) is an integral ideal and \( m_\infty \) is a set of real embeddings of \( F \). We write

\[
(Z_F/m)^\times = (Z_F/m_f)^\times \times \prod_{\sigma \in m_\infty} \{\pm 1\}.
\]

A Hecke character \( \chi \) is said to have modulus \( m \) if \( \chi \) is trivial on the group \( U(m) \) of idèles congruent to 1 mod \( m \):

\[
U(m) = \prod_{p|m_f} (1 + p^{\nu_p(m)}Z_p) \times \prod_{p|m_f} Z_p^\times \times \prod_{\sigma \in m_\infty} \{1\} \times \prod_{\sigma \in m_\infty \setminus \text{real}} \{\pm 1\} \times \prod_{\sigma \text{ complex}} \{1\}.
\]

Equivalently, the conductor of \( \chi \) divides \( m_f \) and \( \chi \) is unramified at all the real places not dividing \( m_\infty \).

The character group of modulus \( m \) is the dual of

\[
C_m = \hat{A}_F^\times / (F^\times \cdot U(m)),
\]

and we have

\[
\hat{C}_F = \bigcup_m \hat{C}_m.
\]

In the remainder of this section, we fix a modulus \( m \).

### 3.2 Explicit description

The character group \( \hat{C}_m \) is isomorphic to \( T^m \times \mathbb{Z}^{n-1} \times \mathbb{R} \) where \( T^m \) is finite. Our goal in the next paragraphs is to prove the following

**Proposition 5.** There exist an integer \( \ell \geq 0 \), a lattice \( \Lambda \) of rank \( \ell + n - 1 \), and two isomorphisms

\[
\mathcal{L} : C_m \xrightarrow{\sim} (\mathbb{Z}^\ell \times \mathbb{R}^n) / \Lambda
\]

\[
\mathcal{L}^* : \hat{C}_m \xrightarrow{\sim} \Lambda^\perp / \mathbb{Z}^\ell
\]

where \( \Lambda^\perp \) is the Pontryagin orthogonal of \( \Lambda \) in \( \mathbb{R}^{\ell+n} \), and such that for all \( \chi \in \hat{C}_m \) and \( x \in \mathbb{A}_F^\times \) we have

\[
\chi(x) = \exp(2i\pi \mathcal{L}^*(\chi) \cdot \mathcal{L}(x)).
\]

The lattice \( \Lambda \) and the isomorphisms \( \mathcal{L} \) and \( \mathcal{L}^* \) will be made explicit in the next subsections.
3.3 Idèle class groups

Definition 6. Let $x \in \mathbb{A}_F^\times$. We define the ideal attached to $x$ to be
\[ \prod_p p^{v_p(x)} \, . \]
Let $S$ be a finite set of primes of $F$. Define the group of $S$-idèles to be
\[ U_S = \prod_{p \notin S} F_p^\times \times \prod_{p \in S} \mathbb{Z}_p^\times \times F_{\mathbb{R}}^\times , \]
and the group of $S$-units $\mathbb{Z}_{F,S}^\times = F^\times \cap U_S$.

Lemma 7. Let $x \in \mathbb{A}_F^\times$. Then $x \in U_S$ if and only if the ideal attached to $x$ belongs to the group $\langle S \rangle$ generated by $S$. If $S$ generates the class group of $F$, then $\mathbb{A}_F^\times = U_S \cdot F^\times$.

Proof. The first property follows from rewriting the definition of $U_S$ as $U_S = \{ x \in \mathbb{A}_F^\times \mid v_p(x) = 0 \text{ for all } p \notin S \}$. Let $x \in \mathbb{A}_F^\times$ and $a$ the ideal attached to $x$. Assuming $S$ generates the class group, let $\alpha$ be such that $a(\alpha^{-1}) \in \langle S \rangle$. Then $x\alpha^{-1} \in U_S$.

Definition 8. Let $S$ be a set of primes generating the class group of $F$. Let
\[ D_S : U_S \to \mathbb{Z}^S \times (\mathbb{Z}_F/m)^\times \times (F_{\mathbb{R}}^\times)^\circ \]
be defined by
\[ D_S(x) = (v_p(x), (u \mod m), (\text{sgn}(x_\sigma))_{\sigma \in \text{m}_F^\infty}, (|x_\sigma|_\sigma \text{ real}, (x_\sigma)_\sigma \text{ complex}), \]
where $u \in \prod_p \mathbb{Z}_p^\times$ is defined by $u_p = x_p \pi_p^{-v_p(x_p)}$, where we recall that $\pi_p \in \mathbb{Z}_p$ is a chosen uniformiser.

Let
\[ D : \mathbb{A}_F^\times / F^\times \to \left[ \mathbb{Z}^S \times (\mathbb{Z}_F/m)^\times \times (F_{\mathbb{R}}^\times)^\circ \right] / D_S(\mathbb{Z}_{F,S}^\times) \]
be defined by $D(x \cdot F^\times) = D_S(x\alpha^{-1})$ where $\alpha \in F^\times$ such that $x\alpha^{-1} \in U_S$.

Lemma 9. Let $S$ be a finite set of primes generating the class group. Then $D$ is well-defined and induces an isomorphism
\[ C_m \cong \left[ \mathbb{Z}^S \times (\mathbb{Z}_F/m)^\times \times (F_{\mathbb{R}}^\times)^\circ \right] / D_S(\mathbb{Z}_{F,S}^\times) . \quad (3) \]

Proof. The existence of the element $\alpha$ from the definition of $D$ exists by Lemma 7. If $x\alpha^{-1}$ and $x\beta^{-1}$ belong to $U_S$ with $\alpha, \beta \in F^\times$, then $\beta/\alpha \in F^\times \cap U_S = \mathbb{Z}_{F,S}^\times$, so $D$ is well-defined. By the decompositions $F_p^\times \cong \pi_p^\times \times \mathbb{Z}_p^\times$ and $F_{\mathbb{R}}^\times \cong \{ \pm 1 \}^{\gamma_1} \times (F_{\mathbb{R}}^\times)^\circ$ and the Chinese remainder theorem, the map $D_S$ is onto, and $\ker D_S = U(m) \subset U_S$. Moreover by definition $D(F^\times) = 1$. This proves that $\ker D = F^\times \cdot U(m)$ and therefore $D$ induces an isomorphism from $C_m = \mathbb{A}_F^\times / (F^\times \cdot U(m))$ to its codomain. \qed
3.4 Logarithm maps

In this section we fix a finite set \( S \) of primes that generates the class group of \( F \) and a modulus \( m \).

**Definition 10.** Consider the usual archimedean logarithm \( \log_\infty : (F_\infty^\times)^o \to \mathbb{R}^{r_1+r_2} \times (\mathbb{R}/\mathbb{Z})^{r_2} = \mathbb{R}^n / \mathbb{Z}^{r_2} \)

\[
\log_\infty(z) = \left( \frac{n_\sigma}{2\pi} \log |z_\sigma|, \frac{\arg(z_\sigma)}{2\pi} \right)_{\sigma \text{ complex}},
\]

(4)

and choose an integer \( r(m) \geq 0 \), a full sublattice \( \Lambda_m \subset \mathbb{Z}^{r(m)} \) and an isomorphism

\[
\log_m : (\mathbb{Z}_F/m)^\times \sim \to \mathbb{Z}^{r(m)}/\Lambda_m.
\]

(5)

Let \( \ell = \#S + r(m) \), and let

\[
\mathcal{L}_S : U_S \to \mathbb{Z}^\ell \times \mathbb{R}^n / \Lambda_m + \mathbb{Z}^{r_2}
\]

be the composition of \( \mathcal{D}_S \) with

\[
\text{Id}_{\mathbb{Z}^S} \times \log_m \times \log_\infty.
\]

We identify \( \Lambda_m \) and \( \mathbb{Z}^{r_2} \) with their embedding in \( \mathbb{Z}^\ell \times \mathbb{R}^n \).

Let

\[
\Lambda = \mathcal{L}_S(\mathbb{Z}_{F,S}^\times) + \Lambda_m + \mathbb{Z}^{r_2},
\]

(6)

and let

\[
\mathcal{L} : \mathbb{K}_F^\times / F^\times \to \mathbb{Z}^\ell \times \mathbb{R}^n / \Lambda
\]

be defined by \( \mathcal{L}(x : F^\times) = \mathcal{L}_S(x \alpha^{-1}) \) where \( \alpha \in F^\times \) is such that \( x \alpha^{-1} \in U_S \).

**Definition 11.** We define the dual logarithm \( \mathcal{L}^* : \hat{C}_m \to (\mathbb{R}/\mathbb{Z})^\ell \times \mathbb{R}^n \) by

\[
\mathcal{L}^*(\chi) = \left( \frac{\arg \chi(p)}{2\pi} \right)_{p \in S}, \left( \frac{\arg \chi(\log^{-1}_m(g_i))}{2\pi} \right)_{i=1}^{r(m)} \left( \varphi_\sigma, k_\sigma \right)_{\sigma \text{ complex}}
\]

(7)

where \( (g_i)_{i=1}^{r(m)} \) is the image in \( \mathbb{Z}^{r(m)}/\Lambda_m \) of the standard basis of \( \mathbb{Z}^{r(m)} \) and \( \varphi_\sigma, k_\sigma \) are the parameters at infinity of \( \chi \).

Recall that we defined \( \chi(p) = \chi(p_\sigma) \), so that \( \mathcal{L}^* \) depends on the choices of \( \pi_\sigma \) for \( p \in S \).

We now prove Proposition [5] in the following precise form.

**Proposition 12.** Let \( \Lambda^\perp \) be the Pontryagin orthogonal of \( \Lambda \) in \( \mathbb{R}^{\ell+n} \). The homomorphisms \( \mathcal{L} \) and \( \mathcal{L}^* \) induce isomorphisms

\[
\mathcal{L} : C_m \to \mathbb{Z}^\ell \times \mathbb{R}^n / \Lambda \text{ and } \mathcal{L}^* : \hat{C}_m \to \Lambda^\perp / \mathbb{Z}^\ell.
\]
Let $\chi \in \hat{C}_m$ be a character of modulus $m$ and let $x \in \mathbb{A}_F^\times$, then

$$\chi(x) = \exp(2i\pi \mathcal{L}^*(\chi) \cdot \mathcal{L}(x)), \quad (8)$$

where $(w, v) \mapsto w \cdot v$ denotes the standard inner product on $\mathbb{R}^{\ell+n}$.

Proof. The fact that $\mathcal{L}$ is well-defined and induces an isomorphism follows immediately from Lemma 9. Applying Pontryagin duality to the sequence

$$0 \to C_m \to \mathbb{R}^{\ell+n}/\Lambda \to (\mathbb{R}/\mathbb{Z})^\ell \to 0$$

gives $\hat{C}_m = \Lambda^\perp/\mathbb{Z}^\ell$.

Let $x \in \mathbb{A}_F^\times$ and write $x = \alpha \cdot x_{\alpha}^{-1}$ with $\alpha \in F^\times$ and $x_{\alpha}^{-1} \in U_S$ by Lemma 7, and let $u$ be as in Definition 8. We have

$$x = \alpha \prod_{p \in S} \pi_p(x_{\alpha}^{-1}) \cdot u \cdot \prod_{\sigma}(x_{\sigma}(\alpha)^{-1}),$$

and therefore

$$\chi(x) = \chi(\alpha) \cdot \prod_{p \in S} \chi_p(\pi_p(x_{\alpha}^{-1})) \cdot \chi(u) \cdot \prod_{\sigma} \chi_{\sigma}(x_{\sigma}(\alpha)^{-1}),$$

where $\chi(\alpha) = 1$ and $\chi(u) = \prod_{p|m} \chi_p(u_p \mod p^m)$. By definition the product of local character evaluations is $\exp(2i\pi \mathcal{L}^*(\chi) \cdot \mathcal{L}(x))$. This also proves that the image of $\mathcal{L}^*$ lies in $\Lambda^\perp$ and that $\mathcal{L}^*$ induces an isomorphism as claimed. \hfill \Box

Remark 13. The lattice $\Lambda$ is not cocompact in $\mathbb{R}^{\ell+n}$, so that the Pontryagin orthogonal $\Lambda^\perp$ is not discrete. In the next section we factor out the norm, so that the resulting lattice is cocompact and its Pontryagin orthogonal can be expressed as a dual lattice as in Section 2.1.

3.5 Characters modulo the norm

Let $C^1_m = C^1_F \cap C_m = \ker(C_m \to \mathbb{R}_{>0})$ be the kernel of the norm, which is compact. We have a canonical splitting inherited from (1)

$$C_m \cong C^1_m \times \mathbb{R}_{>0},$$

and the corresponding decomposition

$$\hat{C}_m \cong \hat{C}^1_m \times \| \cdot \|^\mathbb{R}$$

where $\hat{C}^1_m$ is a discrete finitely generated abelian group.
Proposition 14. Let $v_0 \in \mathbb{R}^{\ell+n}$ be the vector having coordinate $n_\sigma$ at the components corresponding to $\varphi_\sigma$ and 0 elsewhere, and $p_0 : \mathbb{R}^{\ell+n} \rightarrow (\mathbb{R}v_0)^\perp$ the orthogonal projection.

Then $p_0 \circ L$ induces an isomorphism

$$\hat{C}_m^1 \cong p_0(\Lambda)^\vee / \mathbb{Z}^\ell.$$ 

Proof. Let $H = (\mathbb{R}v_0)^\perp = \{ x \mid \sum n_\sigma x_\sigma = 0 \}$, we have an exact sequence

$$0 \rightarrow C_1^m \rightarrow H/p_0(\Lambda) \rightarrow (\mathbb{R}/\mathbb{Z})^\ell \rightarrow 0,$$

where $p_0(\Lambda)$ has full rank in $H$, so that we identify $p_0(\Lambda)^\perp = p_0(\Lambda)^\vee$ in the dual sequence. \hfill \Box

Remark 15. By an appropriate choice of basis of the lattice $\Lambda$, we naturally obtain a structured basis of $C_m$ according to the filtration

$$\hat{\text{Cl}}(m) \subset \hat{C}_m^1 \subset \hat{C}_m.$$ 

It is even possible to obtain a basis exhibiting the filtration

$$\hat{\text{Cl}}_F \subset \hat{\text{Cl}}(m) \subset (\hat{C}_m^1)_{k=0} \subset \hat{C}_m^1 \subset \hat{C}_m,$$

but our implementation makes a different choice of basis, using an SNF basis for the torsion subgroup and exhibiting the subgroup of almost-algebraic characters, as explained in Section 4.

3.6 Algorithms

Since a precise discussion of the complexity is not the main point of the paper, we delegate the difficult operations to oracles.

Definition 16. Let $F$ be a number field and $I_F$ the set of fractional ideals of $\mathbb{Z}_F$. We say that $F$ is strongly computable if it is equipped with

- algorithms to compute field operations in $F$, factorizations into prime ideals and valuations in $I_F$;
- a finite set $S$ of prime ideals generating the class group;
- generators of the $S$-units $\mathbb{Z}_F^S$;
- a principalization oracle $p_S : I_F \rightarrow F^\times \times \mathbb{Z}^S$ such that for every ideal $a \in I_F$ the output $p_S(a) = (\alpha, (a_p)_{p \in S})$ satisfies $a = (\alpha) \prod_{p \in S} p^{a_p}$;
- for each modulus $m$, a lattice $\Lambda_m$ of rank $r(m)$ and a logarithm oracle $\log_m : \mathbb{Z}_F \rightarrow \mathbb{Z}^{r(m)}$ inducing an isomorphism $(\mathbb{Z}_F/m)^\times \cong \mathbb{Z}^{r(m)}/\Lambda_m$. 

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Note that these oracles are available in Pari/GP, using the algorithms described in [4], [6], [7, Section 4.2] and [16].

Using the notations introduced in Definition 10 and Proposition 14 our algorithms are the following.

Algorithm 17.

- Input: a strongly computable number field \( F \), a modulus \( m \) and an ideal \( a \in I_F \).
- Output: a vector \( z \) in \( \mathbb{R}^{\ell+n} \) such that \( L(a) \equiv z \mod \Lambda \).
  1. Let \( (\alpha, (a_p)_p) = p_S(a) \).
  2. Let \( u \in \mathbb{Z}_F \) such that \( \alpha \pi_p^{-v_p(\alpha)} \equiv u \mod p^{v_p(m)} \) for all \( p \mid m_f \).
  3. Return \( z = ((a_p)_{p \in S}, -\log_{m}(u), -\log_{\infty}(\alpha)) \).

Algorithm 18.

- Input: a strongly computable number field \( F \) and a modulus \( m \).
- Output: a matrix \( B \) whose rows generate \( \hat{C}^1_m \) in \( \mathbb{R}^{\ell+n} \).
  1. Let \( A \) be a matrix whose columns form a basis of \( L_S(Z_{F,S}^\chi) + \Lambda_m + \mathbb{Z}^2 + \mathbb{Z}v_0 \) in \( \mathbb{R}^{\ell+n} \).
  2. Let \( B = A^{-1} \): the rows of \( B \) form the basis dual to the columns of \( A \).
  3. Delete from \( B \) the row corresponding to the linear form dual to \( v_0 \).
  4. Replace the rows of \( B \) by their orthogonal projections onto \( (\mathbb{R}v_0)^\perp \).
  5. Return the \((\ell + n - 1) \times (\ell + n)\) matrix \( B \).

Remark 19. These algorithms output numerical approximations in \( \mathbb{R}^{\ell+n} \); their validity to any prescribed accuracy can be certified as follows. In both cases, the numerical approximations come from log embeddings of number field elements, which can be obtained to arbitrary accuracy in polynomial time. All subsequent numerical operations come from linear algebra and can be implemented using certified numerical algorithms [15] with automatic precision increase until a target precision is reached. Our package implements this strategy except that we rely on Pari/GP’s arithmetic which is not certified.

Theorem 20. Algorithm [18] and Algorithm [17] are correct. They are polynomial time, meaning a polynomial number of calls to the oracles with polynomial size input and a polynomial number of other operations.
Proof. Algorithm 18 is correct by Proposition 14. We verify that the value $z$ computed in Algorithm 17 equals $L(a) \mod \Lambda$:

$$\text{let } x = (\pi_p v_p(a)) \text{ be an idèle defining } a, \text{ we have } L(a) = L(x) = L_S(xa^{-1}) \text{ by definition of } L. \text{ Now we have } v_p(xa^{-1}) = a_p \text{ for } p \in S \text{ by definition of } p_S, \text{ and } xa^{-1} \equiv u^{-1} \mod m \text{ by definition of } u. \text{ At infinite places } (xa^{-1})_{\sigma} = a_{\sigma^{-1}}.$$ Hence $L_S(xa^{-1}) \equiv z \mod \Lambda$, and Algorithm 17 is correct. All operations not provided by the oracles can clearly be performed in polynomial time.

4 The subgroup of algebraic characters

Among Hecke quasi-characters, we would like to exhibit the subgroup of algebraic Hecke characters. By Lemma 4, it is equivalent to compute the subgroup of almost-algebraic characters inside the group of Hecke characters. More precisely, let $H_0^\perp \subset \mathbb{R}^{\ell+n}$ be the subgroup of characters defined by $H_0^\perp = \{ \varphi_{\sigma} = 0 \text{ for all } \sigma \}$, then

$$\left(\hat{C}_m\right)^{a.a.} = \hat{C}_m \cap \left(\hat{C}_F\right)^{a.a.} \cong \Lambda^\perp \cap H_0^\perp / \mathbb{Z}^\ell = \left\{ \lambda \in \Lambda^\perp \mid \lambda(h) = 1 \text{ for all } h \in H_0 \right\} / \mathbb{Z}^\ell.$$ However, we do not want to solve the equation $\varphi_{\sigma} = 0$ since the components $\varphi_{\sigma}$ on $\Lambda^\perp$ are only known approximately. We are therefore going to use the known structure of algebraic characters.

Recall that a number field $K$ is CM if it is a totally complex quadratic extension of a totally real field, denoted $K^+$. In this case, the automorphism corresponding to this quadratic extension induces complex conjugation on every complex embedding of $K$, and we therefore denote it by $x \mapsto \bar{x}$.

A classical theorem of Weil and Artin states the following [43, 32]:

- If $F$ does not admit a CM subfield, then every algebraic Hecke character is a finite order character times an integral power of the norm.
- If $F$ admits a CM subfield, then it admits a maximal CM subfield $K$.
  The type of every algebraic Hecke character of $F$ is the lift of the type of an algebraic character of $K$. Equivalently, every almost-algebraic Hecke character of $F$, up to a finite order character, factors through the norm $N_{F/K}$ to $K$.

4.1 Determining the subgroup of algebraic characters from the maximal CM subfield

In this section, we assume that $F$ contains a CM subfield. In particular, $F$ is totally complex.

Let $G = \mathbb{R}^{\ell+n}$ be equipped with its standard inner product and $\Lambda_0 = \Lambda + \mathbb{Z}v_0 = L_S(\mathbb{Z}^{\times}_{F:S}) + \Lambda_m + \mathbb{Z}^2 + \mathbb{Z}v_0$, so that $\hat{C}_m \times \| \cdot \| \mathbb{Z}^\ell \cong \Lambda_0^\perp / \mathbb{Z}^\ell$, with $\Lambda_0^\perp = \Lambda_0^\perp$ in $G$. 

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Our strategy is to capture the algebraic characters in a smaller subspace $H^\perp \subset G$ by using the additional known constraints on almost-algebraic characters, in order to apply the following lemma.

**Lemma 21.** Let $G$ be a finite dimensional $\mathbb{R}$-vector space, let $H \subset G$ be an $\mathbb{R}$-vector subspace and let $\Lambda_0 \subset G$ be a lattice such that $H \cap \Lambda_0$ has full rank in $H$. Then

$$\Lambda_0^\perp \cap H^\perp = \left\{ \lambda \in \Lambda_0^\perp : \lambda \cdot h = 0 \text{ for all } h \in H \cap \Lambda_0 \right\}.$$  

**Proof.** We use the fact that $H$ is an $\mathbb{R}$-subspace generated by $H \cap \Lambda_0$ to write

$$\Lambda_0^\perp \cap H^\perp = \left\{ x \in \Lambda_0^\perp, x \cdot h \in \mathbb{Z} \text{ for all } h \in H \right\}$$

$$= \left\{ x \in \Lambda_0^\perp, x \cdot h = 0 \text{ for all } h \in H \right\}$$

$$= \left\{ x \in \Lambda_0^\perp, x \cdot h = 0 \text{ for all } h \in H \cap \Lambda_0 \right\},$$

proving the claim. \hfill \Box

**Remark 22.** The point of Lemma [21] is that since the inner products between elements of $\Lambda_0$ and $\Lambda_0^\perp$ are in $\mathbb{Z}$, the given expression for $\Lambda_0^\perp \cap H^\perp$ can be computed exactly as a subgroup of $\Lambda_0^\perp$ by linear algebra over $\mathbb{Z}$.

**Example 23.** When $H_0^\perp = \{ \varphi_\sigma = 0 \}$ as above, we have $H_0 = \mathbb{R}r^2$. Then $H_0 \cap \Lambda_0$ is $\mathcal{L}_S(\mathbb{Z}_K^\times, \mathbb{Z}v_0)$, which has rank $r_1(K^+) = r_2(K)$. This has full rank in $H_0$ if and only if $K = F$.

This example shows that using $H_0$ is sufficient when $F$ itself is CM. In the general case, we proceed as follows.

**Proposition 24.** Let $K$ be the maximal CM subfield of $F$, let

$$H^\perp = \{ \varphi_\sigma = 0 \text{ for all } \sigma, \text{ and } (k_\sigma)_\sigma \text{ factors through } K \}$$

and $\Lambda_0 = \mathcal{L}_S(\mathbb{Z}_F^\times, \mathbb{Z}m) + \mathbb{Z}r^2 + \mathbb{Z}v_0$. Then

$$\left( \hat{C} \right)^{a.a.} = (\Lambda_0^\perp \cap H^\perp)/\mathbb{Z}.$$  

where $\Lambda_0 \cap H$ has full rank in $H$. More precisely, the group $U \subset \Lambda_0$ generated by $v_0$, the kernel $\ker(N_{F/K} : \mathbb{Z}r^2 \to \mathbb{Z}r^2(K))$ and $\mathcal{L}_S(u)$ for all $u \in \ker(\mathbb{Z}_F^\times \to (\mathbb{Z}_F/m)^\times)$ such that $N_{F/K}(u) \in K^+$, is contained in $H$ and has full rank.  

**Proof.** Almost-algebraic characters are contained in $H^\perp$ since their infinity types factor through $N_{F/K}$, and we have $H = \mathbb{R}r^2 \times \ker(N_{F/K} : \mathbb{R}r^2 \to \mathbb{R}r^2(K))$. The group $U$ described in the Proposition is clearly contained in $H \cap \Lambda_0$. The map $N_{F/K} : \mathbb{Z}r^2 \to \mathbb{Z}r^2(K)$ is surjective since every complex place of $K$ extends to a complex place of $F$, so that its kernel has rank $r_2 - r_2(K)$. Finally, the units described form a finite index subgroup of $\mathbb{Z}_F^\times$, so the group $U$ has full rank in $H$. \hfill \Box
4.2 The maximal CM subfield

In this section we reformulate the problem of determining the maximal CM subfield in a way that is suitable for an efficient algorithm. Indeed enumerating all subfields, regardless of the algorithm used, could not lead to a polynomial time algorithm since the number of subfields is not polynomially bounded, as the example of multi-quadratic fields shows. One may consider a pure Galois-theoretic approach, but it is currently not known whether one can compute in polynomial time, given a number field $F$, the Galois group of the Galois closure of $F$ (see [1, 17, 12]). Our method relies on the following Lemma.

Lemma 25. Let $F$ be a number field. For $\varepsilon \in \{\pm\}$, let

$$F^\varepsilon = \{ x \in F \mid \sigma(x) = \varepsilon \bar{\sigma}(x) \text{ for all } \sigma \in \text{Hom}(F, \mathbb{C}) \}.$$ 

The following are equivalent:

(i) $F$ admits a CM subfield;

(ii) $F^\varepsilon \neq 0$;

(iii) $\dim_\mathbb{Q} F^+ = \dim_\mathbb{Q} F^-$. 

If those conditions are satisfied, then the largest CM subfield of $F$ is $F^+ + F^-$; it also equals $\mathbb{Q}(a)$ for every $a \in F^-$ having minimal polynomial of degree $2 \dim_\mathbb{Q} F^-$, and such an element exists.

Proof. First note that $F^+$ is the largest totally real subfield of $F$. It is clear that (i) implies (ii). Since $\dim_\mathbb{Q} F^+ \geq 1$, (iii) implies (ii). Let $a, b \in F^-$ be nonzero; then $a/b \in F^+$ and therefore $F^-$ is a one-dimensional vector space over $F^+$, so (ii) implies (iii). Let $a \in F^-$ be nonzero; then $a^2 \in F^+$ is totally negative, so $F^+(a) = F^+ + F^-$ is a CM subfield of $F$, so that (ii) implies (i).

If the conditions are satisfied, then the maximal CM subfield $K$ of $F$ is a quadratic extension of its totally real subfield $F^+$ containing $F^+ + F^-$, so there is equality as claimed. Let $a \in F^- \subset K$ have minimal polynomial of degree $2 \dim_\mathbb{Q} F^- = [K : \mathbb{Q}]$; then it generates $K$ over $\mathbb{Q}$. For every subfield $L \subset K$, if $F^- \subset L$ then $F^+ \subset L$ by taking ratios, so $K \subset L$ and therefore $L = K$. The set of elements of $F^-$ lying in a proper subfield of $F$ is therefore a finite union of proper subspaces, and is therefore nonempty. \hfill $\square$

It is therefore enough to compute $F^-$. Proposition 26 below gives a general algorithm to solve this type of problem.

Proposition 26. Let $F$ be a number field. Let $\Omega$ be a field of characteristic 0, let $R \subset \text{Hom}(F, \Omega)^2$ be a subset and let $(\lambda_r)_{r \in R} \in \mathbb{Q}^R$ be a family of rational numbers. Define

$$F_{R, \lambda} = \{ x \in F \mid \sigma_1(x) = \lambda_r \sigma_2(x) \text{ for all } r = (\sigma_1, \sigma_2) \in R \}.$$ 

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Write $F \otimes F \cong \prod_{i=1}^{k} L_i$ where each $L_i$ is a field. Let $p_i : F \otimes F \to L_i$ be the projection onto $L_i$. For each $r \in R$, let $i(r) \in \{1, \ldots, k\}$ be the index such that $r$ corresponds to an element of $\text{Hom}(L_i, \Omega)$ under the natural bijection

$$\text{Hom}(F, \Omega)^2 \cong \text{Hom}(F \otimes F, \Omega) \cong \bigcup_{i=1}^{k} \text{Hom}(L_i, \Omega),$$

where the last union is disjoint. Let $f : F \to \bigoplus_{r \in R} L_{i(r)}$ be the $\mathbb{Q}$-linear map defined by

$$f(x)_r = p_{i(r)}(x \otimes 1 - \lambda_r (1 \otimes x)) \text{ for all } r \in R.$$

Then $F_{R, \lambda} = \ker f$.

**Proof.** Let $i \in \{1, \ldots, k\}$ and $\varphi \in \text{Hom}(L_i, \Omega)$ correspond to $(\sigma_1, \sigma_2) \in \text{Hom}(F, \Omega)^2$. Then, for all $x \in F$, we have $\sigma_1(x) = \varphi(p_i(x \otimes 1))$ and $\sigma_2(x) = \varphi(p_i(1 \otimes x))$. Noting that $\varphi$ is injective since $L_i$ is a field, we obtain for every $\lambda \in \mathbb{Q}$ the equivalence

$$\sigma_1(x) = \lambda \sigma_2(x) \iff \varphi(p_i(x \otimes 1 - \lambda(1 \otimes x))) = 0 \iff p_i(x \otimes 1 - \lambda(1 \otimes x)) = 0.$$

This proves the claim.

The advantage of rewriting the equations this way is that instead of having conditions in $\Omega$ (which might be a field in which we cannot compute exactly such as $\Omega = \mathbb{C}$ or $\Omega = \mathbb{Q}_p$), the conditions take place in the number fields $L_i$ and $f$ is a linear map between finite-dimensional $\mathbb{Q}$-vector spaces.

**Remarks 27.**

- There are obvious generalizations to conditions expressed with more than two embeddings, but they become more and more expensive as the number of embeddings increases; eventually one may have to compute the full Galois closure of $F$.

- The application to the maximal CM subfield can be generalized to other natural conditions, such as the maximal real subfield, the maximal subfield fixed by some ramification group, or the maximal subfield in which the residue degree of a certain prime divides a given integer.

- When $\lambda_r = 1$ for all $r \in R$, Proposition 26 expresses the subfields of interest as intersections of *principal subfields* in the terminology of van Hoeij, Klüners and Novocin [14].

**4.3 Algorithms**

Section 4.2 leads to the following algorithm to compute the maximal CM subfield.
Algorithm 28.

• Input: an irreducible monic \( P \in \mathbb{Q}[X] \) representing \( F = \mathbb{Q}[X]/(P(X)) \).

• Output: an element \( a \in F \) such that \( \mathbb{Q}(a) \) is the maximal CM subfield of \( F \), or \( \perp \) if \( F \) does not contain a CM subfield.

1. Let \( P(Y) \equiv \prod_{i} Q_{i}(X,Y) \mod P(X) \) be the irreducible factorization of \( P \) over \( F \).

2. Let \( J \) be the set of indices \( i \) such that there exists a complex root \( \alpha \) of \( P \) such that \( Q_{i}(\alpha, \bar{\alpha}) = 0 \).

3. Let \( V \subset F \) be the \( \mathbb{Q} \)-subspace of \( a(X) \mod P(X) \) such that for all \( i \in J \), \( a(X) + a(Y) \equiv 0 \mod (P(X), Q_{i}(X,Y)) \).

4. If \( V = 0 \), return \( \perp \).

5. Let \( a \in V \) be such that the minimal polynomial of \( a \) has degree \( 2 \dim_{\mathbb{Q}} V \). Return \( a \).

Theorem 29. Algorithm 28 is a deterministic polynomial-time algorithm that, given a number field \( F \), computes the maximal CM subfield of \( F \).

Proof. Algorithm 28 is correct by Lemma 25 and Proposition 26 since \( F \otimes_{\mathbb{Q}} F \cong \mathbb{Q}[X,Y]/(P(X), P(Y)) \). It runs in polynomial time because factorization of polynomials over number fields can be performed in polynomial time [20].

We obtain the following algorithm to compute the group of almost-algebraic characters.

Algorithm 30.

• Input: a strongly computable number field \( F \) and a modulus \( m \).

• Output: the group of almost-algebraic characters of modulus \( m \).

1. Let \( K \) be the maximal CM subfield of \( F \), as computed by Algorithm 28.

2. If \( K = \perp \), return the group of finite order characters.

3. Let \( A, B \) be the matrices computed by Algorithm 18 with input \((F, m)\).

4. Let \( U \) be the subgroup described in Proposition 24.

5. Let \( C \) be the subgroup of the row span of \( B \), consisting of elements \( c \) such that \( u \cdot c = 0 \) for all \( u \in U \).

6. Output \( C \).
Theorem 31. Algorithm 30 is correct. It is polynomial time, meaning a polynomial number of calls to the oracles with polynomial size input and a polynomial number of other operations.

Proof. If $F$ does not contain a CM subfield, then almost-algebraic characters are exactly finite order characters by the Artin–Weil theorem. The group $U$ can be computed by linear algebra using the oracles. The group $C$ can be computed by linear algebra over $\mathbb{Q}$ since all the inner products that occur are in $\mathbb{Z}$. The group $C$ is the correct output by the Artin–Weil theorem and Lemma 21 in combination with Proposition 24. All operations not provided by the oracles or Theorem 29 can clearly be performed in polynomial time.

5 Examples

We illustrate the interface of our Pari/GP package \cite{27} with a list of examples of mathematical interest.

5.1 Pari/GP interface

The \texttt{gcharinit}($F$, $m$) function initializes a group structure $\text{gc}$ for a number field $F$ and a modulus $m$. The character group structure $\hat{C}_m \cong \prod_{i=1}^{k} \mathbb{Z}/c_i \mathbb{Z} \times \mathbb{Z}^{n-1} \times \mathbb{R}$ is obtained via the vector $\text{gc.cyc} = [c_1, \ldots, c_k, 0, \ldots, 0, 0]$.

As an example,

```javascript
> gc = gcharinit(x^2+23,3);
> gc.cyc
[6, 0, 0.E-57]
```

expresses the group of Hecke quasi-characters of modulus $m = (3)$ over $F = \mathbb{Q}(\sqrt{-23})$ (see also Equation (1))

$$
\text{Hom}_{\text{cont}}(C_3, \mathbb{C}^\times) = \chi_3^{\mathbb{Z}/6\mathbb{Z}} \times \chi_{C_M} \times \parallel \cdot \parallel_C,
$$

where $\chi_3$ is a character of the ray class group $\text{Cl}_F(3)$ and $\chi_{CM}$ is an infinite order almost-algebraic character.

Characters are described as columns of coordinates in this basis.

```javascript
> gchareval(gc,[1,0,0]~ ,idealprimedec(gc.nf,3)[1])
-0.5000 - 0.8660*I \ \ \text{the prime above 3 is not principal}
> gcharconductor(gc,[2,0,0]~)
[1, []] \ \ \text{a class group character}
```

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The maps $L$ and $L^*$ are accessible as `gcharlog` and `gcharualog`, except that these functions have an extra component corresponding to the norm. For example the character $\chi_{CM}$ has the following parameters in $((\mathbb{R}/\mathbb{Z})^3 \times \mathbb{R} \times \mathbb{Z}) \times \mathbb{C}$, where:

- the set $S$ is $\{(2, \sqrt{-23}-1)\}$;
- the map $\log_m : (\mathbb{Z}_F/m)^\times \to \mathbb{Z}^2/2\mathbb{Z}^2$ is characterized by $\log_m(2) = (1,1)$ and $\log_m(\sqrt{-23}) = (1,0)$.

For closer scrutiny we retrieve the local quasi-characters of $\chi = \chi_{CM}\| \cdot \|$. In particular for a prime $p_3$ dividing the conductor $m = 3$ we obtain a character of the `idealstar` structure $(\mathbb{Z}_F/p_3)^\times$ in addition to a value $\theta \in \mathbb{C}$ such that $\chi(p_3) = \exp(2\pi i \theta)$.

The interface gives a basis of the subgroup of algebraic characters. We can work with these characters via their type.

The $L$-function machinery is readily accessible.

```plaintext
> gcharualog(gc,[0,1,0]~)
[0.11298866677205092301511538301498585720, 0, 1/2, 0, 1, 0]
```

For closer scrutiny we retrieve the local quasi-characters of $\chi = \chi_{CM}\| \cdot \|$. In particular for a prime $p_3$ dividing the conductor $m = 3$ we obtain a character of the `idealstar` structure $(\mathbb{Z}_F/p_3)^\times$ in addition to a value $\theta \in \mathbb{C}$ such that $\chi(p_3) = \exp(2\pi i \theta)$.

The interface gives a basis of the subgroup of algebraic characters. We can work with these characters via their type.

```plaintext
> gcharlocal(gc,[0,1,1]~,1) \ complex place
[1, -I] \ k = 1, phi = -I
> gcharlocal(gc,[0,1,1]~,idealprimedec(gc.nf,3)[2],&grp)
[1, 0.1042940216...+ 0.1748495762...*I] \ [grp char, theta]
> grp.cyc
[2] \ structure of (ZF/p3)^*
```

The $L$-function machinery is readily accessible.
5.2 Small degree examples

We describe explicitly the form taken by infinite order Hecke characters and our choice of basis for low degree fields.

We denote $z = (z_1, \ldots, z_{r_1+r_2})$ the elements of $F_R \cong \mathbb{R}^{r_1} \times \mathbb{C}^{r_2}$. The characters of $(F_R^\times)^o$ are of the form

$$\chi_\infty(z) = \prod_{j=1}^{r_1+r_2} |z_j|^{m_j} \prod_{j=r_1+1}^{r_1+r_2} \left( \frac{z_j}{|z_j|} \right)^{k_j}$$

and conversely, such a character $\chi_\infty$ can be extended to a global Hecke character if it is trivial on a finite index subgroup of $\mathbb{Z}_F^\times$.

Working modulo the norm, we therefore consider the characters of $G_\infty^1 = (F_R^\times)^o / (\mathbb{Z}_F^\times \cdot \mathbb{R}_{>0})$

where $\mathbb{R}_{>0}$ is embedded diagonally. The group $G_\infty^1$ is free of rank $n-1$ and is a full rank lattice in the $\mathbb{Q}$-vector space of all possible parameters at infinity.

When $F$ has class number one and totally positive fundamental units, $G_\infty^1$ is precisely the lattice of infinite order characters of modulus $m = 1$.

**Example 32.** For $F = \mathbb{Q}$, infinite order characters are powers of the norm, and finite order characters are Dirichlet characters.

**Example 33** (real quadratic). Let $F = \mathbb{Q}(\sqrt{D})$ be real quadratic with fundamental unit $\eta_1 > 1$ and regulator $R_F = \log(\eta_1)$. Then $G_\infty^1$ is generated by

$$\chi(z) = \left| \frac{z_1}{z_2} \right|^{\frac{R_F}{2}}.$$

**Example 34** (imaginary quadratic). Let $F = \mathbb{Q}(\sqrt{D})$ be imaginary quadratic with torsion units of order $m$. Then $G_\infty^1$ is generated by

$$\chi(z) = \left( \frac{z_1}{|z_1|} \right)^m.$$

**Example 35** (complex cubic). Let $F$ be complex cubic, and consider a fundamental unit whose complex embeddings are $e^{-\frac{R_F}{2} \pm 2\pi i \alpha}$, where $R_F > 0$ is the regulator and $\alpha \in \mathbb{R}/\mathbb{Z}$ is an angle.

Then $G_\infty^1$ is generated by

$$\chi_1(z) = \left| \frac{z_1}{z_2} \right|^{2\pi \frac{R_F}{2}}$$

and

$$\chi_2(z) = \left| \frac{z_1}{z_2} \right|^{-2\pi \frac{4\alpha}{R_F}} \left( \frac{z_2}{|z_2|} \right)^2.$$
Example 36 (real cubic). Let $F$ be real cubic, and $(\pm e^{\alpha_i}), (\pm e^{\beta_i}) \in F_\mathbb{R}$ the embeddings of two fundamental units, so that the regulator is $R_F = |\alpha_1\beta_2 - \alpha_2\beta_1|$. Then $\hat{G}_\infty^1$ is generated by

$$
\chi_1(z) = |z_1|^{2i\pi \frac{\alpha_1 + 2\alpha_2}{3N_F}} |z_2|^{2i\pi \frac{-2\alpha_1 - \alpha_2}{3N_F}} |z_3|^{2i\pi \frac{-\alpha_1 + \alpha_2}{3N_F}}
$$

and

$$
\chi_2(z) = |z_1|^{2i\pi \frac{\beta_1 + 2\beta_2}{3N_F}} |z_2|^{2i\pi \frac{-2\beta_1 - \beta_2}{3N_F}} |z_3|^{2i\pi \frac{-\beta_1 + \beta_2}{3N_F}}.
$$

5.3 Modular forms

By automorphic induction, Hecke characters of an extension $F/K$ are expected to induce automorphic representations of $GL_{[F:K]}$ over $K$. This is known in a number of cases. Here we provide some explicit examples for quadratic fields, where converse theorems prove the existence of a global automorphic form.

5.3.1 Classical forms over $GL_2$

Let $F = \mathbb{Q}(\sqrt{-D})$ be an imaginary quadratic field of discriminant $-D < 0$ and $k > 0$. To an algebraic character $\chi$ of type $(k,0)$ and conductor $m$ we associate the $q$-series

$$
f_\chi(z) = \sum_{(a,m)=1} \chi(a)q^{N(a)}, q = e^{2i\pi z}, \text{Im}(z) > 0
$$

where the sum runs over integral ideals $a$ coprime to $m$.

Theorem 37 (Hecke[13], Weil[45], Shimura[36, 35]). Let $\chi$ be an algebraic character of type $(k,0)$ and conductor $m$ over $F = \mathbb{Q}(\sqrt{-D})$, then

$$
f_\chi \in S_{k+1}(\Gamma_0(N, \psi_F\psi_\chi))
$$

is a newform of weight $k + 1$, level $N = DN_{F/\mathbb{Q}}(m)$ and character $\psi_F\psi_\chi$ where $\psi_F = (-D)$ is the quadratic character of $F$ and $\psi_\chi(a) = a^{-k}\chi((a))$ is the Dirichlet character of modulus $N_{F/\mathbb{Q}}(m)$ attached to $\chi$.

In the other direction, Ribet proved that all CM newforms come from algebraic Hecke characters [33, Theorem 4.5].

Example 38. Consider $F = \mathbb{Q}(\sqrt{-19})$ and $m = 3$. Our implementation show that up to integral powers of the norm, the algebraic characters are of the form $\chi_3 \chi_\infty^k$ where $\chi_3$ has order 4 and generates $\mathcal{O}(m)$, and $\chi_\infty$ has type $(1,0)$. In Table 1 we list the first algebraic characters and the corresponding CM modular forms referenced in [21].

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Table 1: Some modular forms with CM by $\mathbb{Q}(\sqrt{-19})$

| $(i,k)$ | quasi-character | modular form | first zero |
|---------|----------------|--------------|------------|
| $(1,0)$ | $[1,0,0]$ | 171.1.c.a.37.1 | 2.55662379... |
| $(2,0)$ | $[2,0,0]$ | Dirichlet 57.56 | 2.4031422... |
| $(3,0)$ | $[3,0,0]$ | 171.1.c.a.37.1 | 2.55662379... |
| $(0,1)$ | $[0,-1,-1/2]$ | 171.2.d.a.170.3 | 1.19761556... |
| $(1,1)$ | $[1,-1,-1/2]$ | 171.2.d.a.170.1 | 3.031017177... |
| $(2,1)$ | $[2,-1,-1/2]$ | 171.2.d.a.170.2 | 2.19220898... |
| $(3,1)$ | $[3,-1,-1/2]$ | 171.2.d.a.170.4 | 0.57935987... |
| $(0,2)$ | $[0,-2,-1]$ | 171.3.c.d.37.2 | 1.76815328... |
| $(1,2)$ | $[1,-2,-1]$ | 171.3.c.a.37.1 | 1.8455925... |
| $(2,2)$ | $[2,-2,-1]$ | 171.3.c.d.37.1 | 1.54865425... |
| $(3,2)$ | $[3,-2,-1]$ | 19.3.b.a.18.1 | 3.78194741... |
| $(0,3)$ | $[0,-3,-3/2]$ | 171.4.d.a.170.4 | 1.59003776... |
| $(1,3)$ | $[1,-3,-3/2]$ | 171.4.d.a.170.3 | 1.36085197... |
| $(2,3)$ | $[2,-3,-3/2]$ | 171.4.d.a.170.1 | 0.08123213... |
| $(3,3)$ | $[3,-3,-3/2]$ | 171.4.d.a.170.2 | 0.70404412... |

### 5.3.2 Maass waveforms

Let $F = \mathbb{Q}(\sqrt{D})$ be a real quadratic field of discriminant $D$ and fundamental unit $\eta_1 > 1$, and $\chi_m$ a Hecke character of conductor $m = (\infty_1 \infty_2)^\epsilon$ for $\epsilon \in \{0, 1\}$ whose restriction to $F \times \mathbb{R}$ is

$$\chi_m(z) = \text{sgn}(z_1 z_2) \epsilon |\frac{z_1}{z_2}|^{ir_m}, \quad r_m = \frac{m\pi}{2\log(\eta_1)},$$

where $\epsilon \equiv m \mod 2$.

It corresponds to a CM Maass form [8, section 15.3.10].

**Proposition 39.** Let $\cos^{(0)}(x) = \cos(x)$ and $\cos^{(-1)}(x) = \sin(x)$, and $K_{ir}$ denote the modified Bessel function of the second kind of parameter $ir$. The function

$$f(x + iy) = \sqrt{y} \sum_a \chi_m(a) K_{ir_m}(2\pi N(a)y) \cos^{(-\epsilon)}(2\pi N(a)x)$$

is a cusp form of weight 0 and character $\psi_F$ on $\Gamma_0(D)$ with Laplace eigenvalue $\lambda_m = \frac{1}{4} + r_m^2$, where $\psi_F = (\frac{D}{\cdot})$ is the quadratic character of $F$.

**Example 40.** Let $F = \mathbb{Q}(\sqrt{5})$, this field has trivial class group and fundamental unit $\eta = \frac{1 + \sqrt{5}}{2}$. The character $\chi_m$ above is an actual Hecke character of modulus $m = (\infty_1 \infty_2)^\epsilon$. Using the $L$-function facilities in Pari/GP we compute the first zero $0 < \gamma_1$ such that $L(\chi_m, \frac{1}{2} + i\gamma_1) = 0$. Results are shown in Table 2.
Table 2: First zero of Maass form $L$-functions of real quadratic field $\mathbb{Q}(\sqrt{5})$.

Note that we obtain arbitrary large imaginary spectral parameters: this raises computational issues on the $L$-function side which are currently not addressed in Pari/GP. See [2] for the case of degree 2 Maass forms.

5.4 CM abelian varieties

In this section we give examples of CM abelian varieties and the corresponding algebraic Hecke characters. We insist on proving equalities of $L$-functions rather than observing a numerical coincidence, as this is possible thanks to CM theory. For the general terminology of CM theory, we refer to [19, 25]. The following is a special case of [19, Chapter 4 Theorem 6.2].

**Theorem 41** (Shimura [37], Milne [26]). Let $A/\mathbb{Q}$ be a simple abelian variety of dimension $g$. Let $K$ be a CM field of degree $2g$ and $\iota: K \to \text{End}^0(A)$ an embedding, and let $\Phi$ be the corresponding CM type on $K$. Let $F$ be the field of definition of $\iota(K)$, and let $\Phi^*$ be the dual type on $F$. Then $F/\mathbb{Q}$ is Galois; let $G = \text{Gal}(F/\mathbb{Q})$. Let $\pi$ be the injective morphism $\pi: G \to \text{Aut}(K)$ such that $\iota(\lambda)^{\sigma} = \iota(\lambda^{\pi(\sigma)})$ for all $\lambda \in K$ and $\sigma \in G$. Then there exists an algebraic Hecke character $\chi$ over $F$ of type $\Phi^*$ and valued in $K$ such that

$$L(A, s) = \prod_{\tau \in \text{Hom}(K, \mathbb{C})/\pi(G)} L(\chi^{\tau}, s).$$

**Example 42.** Let $A$ be the Jacobian of the genus 2 curve 28561.a.371293.1 from the LMFDB [22] given by the equation

$$y^2 + x^3 y = -2x^4 - 2x^3 + 2x^2 + 3x - 2.$$
Let \( K = \mathbb{Q}[x]/(x^4 - x^3 + 2x^2 + 4x + 3) = \mathbb{Q}(\alpha) \) be the unique degree 4 subfield of \( \mathbb{Q}(\zeta_{13}) \). The surface \( A \) is simple, has CM by \( K \), and all endomorphisms of \( A \) are defined over \( K \), as recorded in the LMFDB and proved by the algorithms of \cite{23} \cite{24}. We therefore have \( F = K \) in the notation of Theorem \cite{41}

Since \( K/\mathbb{Q} \) is Galois, \( \pi(G) \) acts transitively on \( \text{Hom}(K, \mathbb{C}) \). All CM types of \( K \) are in the same Galois orbit; let \( \Phi^* = \{ \alpha \mapsto -0.65 \ldots + 0.52 \ldots i, \alpha \mapsto 1.15 \ldots + 1.72 \ldots i \} \). By Theorem \cite{41} there exists an algebraic Hecke character \( \chi \) of \( K \) of type \( \Phi^* \) such that

\[
L(A, s) = L(\chi, s).
\]

The conductor of \( A \) is 28561 = 13^4, and the discriminant of \( K \) is 2197 = 13^3. Moreover, \( K \) has a unique prime \( \mathfrak{p} \) above 13, so the conductor of \( \chi \) must be \( \mathfrak{p} \).

Using our implementation we compute the group of characters of modulus \( \mathfrak{p} \). The subgroup of finite order characters has order 3, and there exists an algebraic character, unique up to multiplication by a finite order character, of type \( \Phi^* \). Among the three algebraic characters of this type, two have a non-real \( L \)-function coefficient \( a_3 \), and therefore cannot be \( \chi \). So \( \chi \) is the remaining one, which is uniquely characterized by its type and the approximate value

\[
\chi(q) = -1.65138 \ldots - 0.52241 \ldots i
\]

where \( q = (3, \alpha) \) (label 3.1 as defined in \cite{10}). The restriction of \( \chi \) to \( (\mathbb{Z}_K/\mathfrak{p})^* \) has order 2. The values of \( \chi \) at some prime ideals are given in Table \ref{table:values}

| prime \( \mathfrak{p} \) | \( \chi(\mathfrak{p}) \in \mathbb{C} \) | \( \chi(\mathfrak{p}) \in K \) |
|-------------------|---------------------------------|--------------------------|
| 3.1               | -1.65138 \ldots - 0.52241 \ldots i | -\frac{2}{5} \alpha^3 - \frac{3}{5} \alpha - 2 |
| 3.2               | 0.15138 \ldots - 1.72542 \ldots i | -\frac{1}{3} \alpha^3 + \alpha^2 - \frac{4}{5} \alpha - 1 |
| 3.3               | -1.65138 \ldots + 0.52241 \ldots i | \alpha - 1 |
| 3.4               | 0.15138 \ldots + 1.72542 \ldots i | \frac{2}{5} \alpha^3 - \alpha^2 + \frac{1}{5} \alpha + 1 |
| 13.1              | \pm 3.60555 \ldots | \pm \sqrt{13} = \pm \left( \frac{2}{5} \alpha^3 - \frac{3}{5} \alpha + 3 \right) |
| 16.1              | -4 | -4 |
| 29.1              | -3.45416 \ldots - 4.13143 \ldots i | -\frac{5}{3} \alpha^3 + 3 \alpha^2 - \frac{7}{3} \alpha - 5 |
| 29.2              | 1.95416 \ldots + 5.01809 \ldots i | 2 \alpha^3 - 2 \alpha^2 + 5 \alpha + 5 |
| 29.3              | -3.45416 \ldots + 4.13143 \ldots i | \frac{2}{5} \alpha^3 - 3 \alpha^2 + \frac{10}{3} \alpha - 1 |
| 29.4              | 1.95416 \ldots - 5.01809 \ldots i | -\alpha^3 + 2 \alpha^2 - 6 \alpha - 2 |

Table 3: Values of the algebraic character \( \chi \) attached to an abelian surface

**Example 43.** Let \( A \) be the Jacobian of the genus 3 curve \( C : y^3 = x(x^3 - 1) \) from the LMFDB \cite{23}

\[
C : y^3 = x(x^3 - 1).
\]

Let \( K = \mathbb{Q}(\zeta_9) \). The curve \( C \) has an automorphism of order 9, defined over \( K \) and given by \( (x, y) \mapsto (\zeta_3^3 x, \zeta_9 y) \). In particular, the threefold \( A \) has
CM by $K$ defined over $K$. By point counting, the Euler polynomial of $A$ at $p = 7$ is

$$1 + pT^3 + p^3T^6,$$

which is irreducible over $\mathbb{Q}$, proving that $A$ is simple. Since $K/\mathbb{Q}$ is Galois, $\pi(G)$ acts transitively on $\text{Hom}(K, \mathbb{C})$ in the notations of Theorem 41. There are two Galois orbits of CM types on $K$: one lifted from the CM subfield $\mathbb{Q}(\zeta_3) \subset K$, and a primitive one. Let $\Phi^* = \{\zeta_9 \mapsto \exp(2i\pi \zeta_3^2), \zeta_9 \mapsto \exp(2i\pi \zeta_3^2), \zeta_9 \mapsto \exp(2i\pi \zeta_3^2)\}$, which is primitive. By Theorem 41, there exists an algebraic Hecke character $\chi$ of $K$ of type $\Phi^*$ with values in $K$ such that

$$L(A, s) = L(\chi, s).$$

Let $p$ be the unique prime of $K$ above 3. By computing resultants we see that $A$ has good reduction away from 3. In particular the conductor of $\chi$ is a power of $p$, say $p^m$. The restriction of $\chi$ to $(\mathbb{Z}_K/p^m)^\times$ has finite order and takes values in $K$, and therefore has order dividing 18. By studying the 3-adic convergence of $(1 + x)^{1/18}$ we see that $1 + p^{16} \subset (K_p^\times)^{18}$ and in particular we have $m \leq 16$. Alternatively, we could bound $m$ by using the reduction theory of Picard curves [3], but the above method works in cases where no reduction theory is available.

| prime $\tau$ | $\chi(\tau) \in \mathbb{C}$ | $\chi(\tau) \in K$ |
|-------------|--------------------------|-----------------|
| 3.1         | $\langle -\zeta_9\rangle \sqrt{-3} = \langle -\zeta_9\rangle (1 + 2\zeta_9^3)$ | $\langle -\zeta_9\rangle \sqrt{-3} = \langle -\zeta_9\rangle (1 + 2\zeta_9^3)$ |
| 19.1        | $4.34002 \ldots + 0.40522 \ldots i$ | $2\zeta_9^5 + 2\zeta_9^4 + 2\zeta_9^3 + \zeta_9^2 - 2\zeta_9 + 2$ |
| 19.2        | $-4.11721 \ldots + 1.43128 \ldots i$ | $-\zeta_9^5 + 2\zeta_9^4 + 2\zeta_9^3 - 2\zeta_9^2 + 4\zeta_9 + 2$ |
| 19.3        | $4.34002 \ldots - 0.40522 \ldots i$ | $4\zeta_9^5 + \zeta_9^4 - 2\zeta_9^3 + 2\zeta_9^2 - \zeta_9$ |
| 19.4        | $-4.11721 \ldots - 1.43128 \ldots i$ | $-2\zeta_9^5 + \zeta_9^4 - 2\zeta_9^3 - 2\zeta_9^2 + 4\zeta_9 - \zeta_9$ |
| 19.5        | $2.77718 \ldots + 3.35964 \ldots i$ | $-\zeta_9^5 - 4\zeta_9^4 + 2\zeta_9^3 + \zeta_9^2 - 2\zeta_9 + 2$ |
| 19.6        | $2.77718 \ldots - 3.35964 \ldots i$ | $-2\zeta_9^5 - 2\zeta_9^4 - 2\zeta_9^3 + 2\zeta_9^2 - \zeta_9$ |
| 37.1        | $4.34002 \ldots - 4.26194 \ldots i$ | $4\zeta_9^5 + 4\zeta_9^4 - 2\zeta_9^3 + 5\zeta_9^2 + 2\zeta_9$ |
| 37.2        | $2.77718 \ldots - 5.41176 \ldots i$ | $-5\zeta_9^5 - 2\zeta_9^4 - 2\zeta_9^3 - \zeta_9^2 - 4\zeta_9$ |
| 37.3        | $-4.11721 \ldots - 4.47756 \ldots i$ | $-4\zeta_9^5 + 5\zeta_9^4 + 2\zeta_9^3 - 2\zeta_9^2 + 4\zeta_9 + 2$ |
| 37.4        | $4.34002 \ldots + 4.26194 \ldots i$ | $2\zeta_9^5 - \zeta_9^4 + 2\zeta_9^3 - 2\zeta_9^2 - 5\zeta_9 + 2$ |
| 37.5        | $2.77718 \ldots + 5.41176 \ldots i$ | $2\zeta_9^5 - 4\zeta_9^4 + 2\zeta_9^3 + 4\zeta_9^2 + \zeta_9 + 2$ |
| 37.6        | $-4.11721 \ldots + 4.47756 \ldots i$ | $\zeta_9^5 - 2\zeta_9^4 - 2\zeta_9^3 - 4\zeta_9^2 + 2\zeta_9$ |
| 64.1        | $-8$ | $-8$ |

Table 4: Values of the algebraic character $\chi$ attached to an abelian threefold

Using our implementation we compute the group of characters of modulus $p^{36}$. The subgroup of finite order characters is isomorphic to $C_3^4$. There exists an algebraic character of type $\Phi^*$, unique up to multiplication by a finite order character. Out of these $9^4 = 6561$ candidate characters, checking that the value of $a_{19}$ is sufficiently close to the value for $A$, namely $a_{19}(A) = 6$, eliminates all but 2 candidates. Checking that the value
of \( a_{109} \) is sufficiently close to \( a_{109}(A) = -21 \) leaves only one remaining candidate, which must therefore be \( \chi \). The conductor of \( \chi \) is \( p^4 \) and \( \chi \) is in fact the unique algebraic character of type \( \Phi^* \) and conductor \( p^4 \), and the restriction of \( \chi \) to \( (\mathbb{Z}_K/p^4)^* \) has order 18. The values of \( \chi \) at some prime ideals\(^2\) are given in Table 4.

### 5.5 Density of gamma shifts

The spectral parameters of an \( L \)-function are the gamma shifts \( \mu_j \) appearing in the gamma factor

\[
\gamma(s) = \prod_{j=1}^{r_1} \Gamma_{\mathbb{R}}(s + \mu_j) \prod_{j=r_1+1}^{r_1+r_2} \Gamma_{\mathbb{C}}(s + \mu_j).
\]

of its normalized functional equation \( L(s)\gamma(s) = \Lambda(s) = \varepsilon\Lambda(1-s) \). In this setting, the real parts \( \text{Re}(\mu_j) \leq r_1 \) and \( \text{Re}(2\mu_j) > r_1 \) are expected to be integers, whereas the imaginary parts can be arbitrary transcendentals subject to

\[
\sum_{j=1}^{r_1} \mu_j + \sum_{j=r_1+1}^{r_1+r_2} 2\mu_j \in \mathbb{R}.
\]

As a matter of fact, Hecke characters allow us to attain a dense subspace of these possible gamma shifts. The following statement must be well-known but we could not find a reference for it.

**Proposition 44.** Let \( r_1, r_2 \geq 0 \) and \( (\mu_j^*) \in (\{0,1\} + i\mathbb{R})^{r_1} \times (\frac{1}{2}\mathbb{Z}_{\geq 0} + i\mathbb{R})^{r_2} \) a family of spectral parameters such that \( \sum_{j\leq r_1} \mu_j^* + 2\sum_{j>r_1} \mu_j^* \in \mathbb{R} \).

Then for every number field \( F \) of signature \( (r_1, r_2) \) and every \( \epsilon > 0 \), there exists a Hecke character \( \chi \) of \( F \) whose \( L \)-function gamma shifts \( \mu_j(\chi) \) satisfy

\[
|\mu_j(\chi) - \mu_j^*| < \epsilon.
\]

**Proof.** Let \( F \) be a number field of signature \( (r_1, r_2) \). For every modulus \( m \), let \( G_m \subset \hat{\mathbb{C}} \) be the image of the map \( \hat{C}_m \to \hat{\mathbb{C}} \), that is, the group of infinity-types of characters of modulus \( m \). The group \( G_m \) is the group of elements \( \chi \in \hat{\mathbb{C}} \) such that \( \chi(u) = 1 \) for all \( u \in \mathbb{Z}_F^*(m) = \ker(\mathbb{Z}_F^* \to (\mathbb{Z}/m)^*) \).

Let \( M > 0 \) be an integer. By the congruence subgroup property for unit groups of number fields \[5\] Théorème 1], there exists a modulus \( m \) such that \( \mathbb{Z}_F^*(m) \subset (\mathbb{Z}/m)^* \). In particular, we get that

\[
\left\{ \chi \in \hat{\mathbb{C}} \mid \chi^M \in G_1 \right\} \subset G_m.
\]

Since the image of \( G_1 \) in \( \mathbb{R}^{r_1+r_2} \times \mathbb{Z}^{r_2} \) has full rank, this proves that \( \bigcup_m G_m \) is dense in \( \hat{\mathbb{C}} \), which implies the claim. \( \Box \)

\(^2\)Labels are as in \[10\] but with respect to the cyclotomic polynomial \( \Phi_9 \), which is not the \texttt{polredabs} polynomial.
This makes Hecke characters good test cases for \(L\)-functions software, since their coefficients are relatively easy to compute compared to other transcendental automorphic forms.

**Example 45.** We exhibit a character of conductor \(2^{20}\) over the real cubic field \(F = \mathbb{Q}[x]/(x^3 - 3x + 1)\) whose parameters \(\varphi_1\) and \(\varphi_2\) approximate the constants \(\pi\) and \(e\) to 5 digits.

```bash
> g=gcharinit(x^3-3*x+1,2^20); chi = [0,-2033118, 694865]~;
> gcharlocal(g,chi,1)
[0, 3.14159223855113838337575888544915179]
> gcharlocal(g,chi,2)
[0, 2.7182831477529933175766620889117919084]
```

5.6 **Partially algebraic Hecke characters**

In view of the special role played by algebraic Hecke characters, it is natural to ask whether there exists partially algebraic Hecke characters, that is, characters such that \(\varphi_\sigma = 0\) for some \(\sigma\) but not all.\(^3\) We provide a construction of such characters.

**Proposition 46.** Assume \(F\) is a quadratic extension of another number field \(F_0\). Let \(R\) be the set of real places of \(F_0\) that become complex in \(F\), and let \(n_0\) be the degree of \(F_0\). Then for every modulus \(m\) of \(F\), there exists a subgroup \(H\) of \(\hat{\mathbb{C}}_m\) of rank \(n_0\) in which every character satisfies \(\varphi_\sigma = 0\) for every \(\sigma\in R\).

**Proof.** It suffices to prove the statement for the modulus \(m = 1\). Let \(g\) be the nontrivial element of \(\text{Gal}(F/F_0)\), which acts on \(\hat{C}_1\). Let \(H\) be the subgroup of \(\chi \in \hat{C}_1\) such that there exists a finite order \(\xi \in \hat{C}_1\) with \(\chi^g = \xi \chi^{-1}\). We have \(\text{rk}(\hat{C}_1(F)) = n - 1 = 2n_0 - 1\) and \(\text{rk}(\hat{C}_1(F_0)) = n_0 - 1\) (as is well-known but also easily seen from Proposition 14), so the rank of \(H\) is exactly \(n_0\). Moreover, for every infinite place \(\sigma\) of \(F\), every element of \(H\) satisfies \(\varphi_\sigma = -\varphi_\sigma\). In particular for \(\sigma \in R\) this means that \(\varphi_\sigma = 0\).

**Corollary 47.** Under the same hypotheses as Proposition 46, let \(r = 0\) if \(F\) does not contain a CM subfield and \(r\) be the degree of the maximal real subfield of \(F\) otherwise. Then for every modulus \(m\) of \(F\), there exists a subgroup \(H\) of \(\hat{C}_m\) of rank \(n_0 - r\) in which every character satisfies \(\varphi_\sigma = 0\) for every \(\sigma \in R\) and such that \(H\) contains no nonzero almost-algebraic character. In particular, if \(F\) is not CM then there exists a partially algebraic character over \(F\).

\(^3\)See [https://mathoverflow.net/questions/310706](https://mathoverflow.net/questions/310706)
Proof. The integer \( r \) is the rank of the group of almost-algebraic characters. \( \square \)

Example 48. Consider \( F_0 = \mathbb{Q}(\sqrt{5}) \subset F = \mathbb{Q}(5^{1/4}) \).

\begin{verbatim}
> gc=gcharinit(x^4-5,1);
> chi = [1,0,0]~;
> gcharlocal(gc,chi,1)
[0, -0.72908519629282042564585827345932876864]
> gcharlocal(gc,chi,2)
[0, 0.72908519629282042564585827345932876864]
> gcharlocal(gc,chi,3)
[2, 0]
\end{verbatim}

The character \( \chi \) satisfies
\[
\chi_{\sigma_1}: x \mapsto |x|^{0.729...}, \quad \chi_{\sigma_2}: x \mapsto |x|^{0.729...}, \quad \text{and} \quad \chi_{\sigma_3}: z \mapsto (z/|z|)^2,
\]
and is therefore an example of a partially algebraic character. Since \( n_0 = 2 \) there is another independent partially algebraic character (namely \([0, 1, 0]^\sim\)).

In a general number field \( F \), if one fixes a set of infinite places \( \Sigma \), a natural question is to determine the group of \( \Sigma \)-algebraic characters, i.e. characters such that \( \varphi_\sigma = 0 \) for every \( \sigma \in \Sigma \). The field \( F \) contains a maximal subfield \( K_0 \) that is real at places below \( \Sigma \), and may contain a quadratic extension \( K \) of \( K_0 \) in which all places below \( \Sigma \) are complex. When this is the case, one obtains a corresponding group of \( \Sigma \)-algebraic characters. Does this construction account for all the possible infinity types? Unlike the algebraic case where Galois theory is sufficient to obtain a complete characterisation, the general case seems to involve transcendence problems.

By automorphic induction to \( \text{GL}_2 \), partially algebraic characters yield automorphic representations that are non-algebraic principal series at some infinite places and discrete series at other ones. Analogously to \([30]\), one may ask to explicitly construct such "partial Maass forms" that do not come from Hecke characters. A possible way of doing this would be to compute Maass forms on a well-chosen quaternion algebra and to use the Jacquet–Langlands correspondence.

5.7 Twists and special values

Another interesting use of Hecke characters is to twist other \( L \)-functions to obtain new ones. Our implementation makes it easy to follow the experiments of \([40]\) on twists of elliptic curve \( L \)-functions.

Let \( E/F \) be an elliptic curve of conductor \( N_{E/F} \) over an imaginary quadratic field \( F \), and \( \chi \) be an algebraic Hecke character of type \((a,b)\) and conductor \( f \) over \( F \).
Assume gcd(f, N_{E/F}) = 1, then the twist
\[ L(E \otimes \chi, s) = \sum_{(n,f) = 1} a_n(E) \chi(n) N(n)^{-s} \]
conjecturally satisfies the functional equation
\[ \Lambda(E \otimes \chi, s) = W \Lambda(E \otimes \overline{\chi}, 1 + a + b - s) \]
where
\[ \Lambda(E \otimes \chi, s) = (N(f)^2 N_{E/F})^{\frac{s}{2}} \Gamma_C(s - \min(a,b)) \Gamma_C(s - \min(a,b) - 1) L(E \otimes \chi, s) \]
with special values predicted by Deligne’s period conjecture [11].

Example 49. Let \( F = \mathbb{Q}(\sqrt{-43}) \), \( E/F \) the curve \([43.1.a.1]\) of equation \( y^2 + y = x^3 + x^2 \), and \( \chi \) the algebraic character of conductor 1 and type \((-2,2)\).

We check numerically that the special value is a period related to \( F \).
\[ L(E \otimes \chi, 1) \approx 2.996120826544463 \ldots \approx \frac{2\pi}{\sqrt{43}} \Omega_F^8, \text{ where } \Omega_F = \sqrt{\prod_{a=1}^{42} \Gamma\left(\frac{a}{43}\right)^{-1}}. \]

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A Implementation notes

This appendix collects notes on the matrix transformations used in our implementation [27].

Let $\zeta$ be a generator of the group of roots of unity of $F$, let $\Lambda_u$ (computed with $\text{bnfinit.fu}$) be the image by $L_S$ of the span of a basis of $\mathbb{Z}_F^\times/\langle \zeta \rangle$, and let $\Lambda_S$ (computed with $\text{bnfsunit[1]}$) be the image by $L_S$ of the span of a basis of $\mathbb{Z}_{F,S}^\times/\mathbb{Z}_F^\times$. Let $\mathbb{Z}_F^\times(m) = \ker(\mathbb{Z}_F^\times \to (\mathbb{Z}_F/m)^\times)$. We define the following subgroup:

$$(\hat{C}_m^1)_{k=0} = \left\{ \chi \in \hat{C}_m^1 \mid k_\sigma = 0 \text{ for every complex embedding } \sigma \right\}.$$ 

We will describe a sequence of matrices representing a generating set of $\Lambda_0 = \Lambda + \mathbb{Z}v_0 = \Lambda \oplus \mathbb{Z}v_0$. We only write the ring to which the coefficients of the matrices belong. We also indicate the number of rows and columns, with the following notations: $n_s = |S|$, the integer $n_c$ is the rank of $\Lambda_m$ ($r(m)$ in the main paper, but here we follow the notations from the code).

We want to apply matrix operations so that we compute a basis of $\Lambda$ from the generating set, exhibit interesting subgroups of the group of Hecke characters, and preserve exactness of coefficients whenever possible.

At each step, we apply column operations to modify the generating set of $\Lambda$. These column operations are obtained by applying a HNF reduction to the submatrix displayed as a red block. As a matter of fact, we rely on the following property: from two lattices $G, H$ the HNF computes a subgroup $H'$ of $H$ that is saturated (i.e. the intersection of $H$ with a vector space), defined by some rows being $0$, and a complement $G'$ of $H'$ in $G + H$.

At the end we compute an inverse to get a basis of $\Lambda_0'$ from which we deduce a basis of $\Lambda^\vee$, and we describe various subgroups of the group of Hecke characters that appear naturally. In the end tables, the meaning of the rows is as follows: the title of the row is the subgroup generated by all the previous rows. In other words, the corresponding rows generate a complement of the previous rows in the subgroup in the title of the row. The column labelled $\chi(p)$ contains the values $\frac{1}{2\pi} \arg \chi(p)$ and the column labelled $\chi_m(g_i)$ contains the values $\frac{1}{2\pi} \arg \chi(\log_m(g_i))$.

A.1 Case without a CM subfield

Here we do not assume that we have a CM subfield. This subsection is no longer implemented, but serves as a simpler version of the next subsection.
Initial matrix.

\[
\begin{array}{cccccc}
\text{n_s} & r_1 + r_2 - 1 & \text{n_c} & 1 & \text{r_2} \\
\Lambda_S & \Lambda_u & \Lambda_m & \zeta & \mathbb{Z}^2 \\
\end{array}
\]

| n_s | v_S | n_c | \log_m | r_1 + r_2 | \arg_\sigma |
|-----|------|-----|--------|-----------|------------|
| n_s | Z | 0 | 0 | 0 | 0 |
| n_c | Z | Z | Z | Z | 0 |
| r_1 + r_2 | \log_\sigma | R | R | 0 | 0 |
| r_2 | \arg_\sigma | R | R | 0 | \mathbb{Q} | \mathbb{Z} |

Step 1. We compute the subgroup \( \langle \zeta(m) \rangle = \langle \zeta \rangle \cap \mathbb{Z}_F^\times(m) \) and a complement \( \Lambda_{m,\zeta} \) of \( \langle \zeta(m) \rangle \) in \( \Lambda_m + \langle \zeta \rangle \).

| n_s | r_1 + r_2 - 1 | \text{n_c} | 1 | \text{r_2} \\
|-----|----------------|-----|-----|-----|
| \Lambda_S & \Lambda_u & \Lambda_{m,\zeta} & \langle \zeta(m) \rangle & \mathbb{Z}^2 |
| n_s | v_S | n_c | \log_m | r_1 + r_2 | \arg_\sigma |
|-----|------|-----|--------|-----------|------------|
| n_s | Z | 0 | 0 | 0 | 0 |
| n_c | Z | Z | Z | Z | 0 |
| r_1 + r_2 | \log_\sigma | R | R | 0 | 0 |
| r_2 | \arg_\sigma | R | R | 0 | \mathbb{Q} | \mathbb{Z} |

Step 2. We compute the span \( \Lambda_u(m) \subset \Lambda_u \) of a basis of \( \mathbb{Z}_F^\times(m)/\langle \zeta(m) \rangle \) and a complement \( \Lambda_{m,u} \) of \( \Lambda_u(m) \) in \( \Lambda_{m,\zeta} + \Lambda_u \).

| n_s | n_c | r_1 + r_2 - 1 | \text{r_2} |
|-----|-----|----------------|-----|
| \Lambda_S & \Lambda_{m,u} & \Lambda_u(m) & \langle \zeta(m) \rangle & \mathbb{Z}^2 |
| n_s | v_S | n_c | \log_m | r_1 + r_2 | \arg_\sigma |
|-----|------|-----|--------|-----------|------------|
| n_s | Z | 0 | 0 | 0 | 0 |
| n_c | Z | Z | Z | Z | 0 |
| r_1 + r_2 | \log_\sigma | R | R | 0 | 0 |
| r_2 | \arg_\sigma | R | R | 0 | \mathbb{Q} | \mathbb{Z} |

Step 3. We compute \( \Lambda_\zeta = \mathbb{Z}^2 + \langle \zeta(m) \rangle \).

| n_s | n_c | r_1 + r_2 - 1 | \text{r_2} |
|-----|-----|----------------|-----|
| \Lambda_S & \Lambda_{m,u} & \Lambda_u(m) & \Lambda_\zeta |
| n_s | v_S | n_c | \log_m | r_1 + r_2 | \arg_\sigma |
|-----|------|-----|--------|-----------|------------|
| n_s | Z | 0 | 0 | 0 | 0 |
| n_c | Z | Z | Z | Z | 0 |
| r_1 + r_2 | \log_\sigma | R | R | 0 | \mathbb{Z} |
| r_2 | \arg_\sigma | R | R | 0 | \mathbb{Q} |

Step 4. We include \( v_0 \) to obtain a square block on the \( \log_\sigma \) components.

| n_s | n_c | r_1 + r_2 - 1 | \text{r_2} |
|-----|-----|----------------|-----|
| \Lambda_S & \Lambda_{m,u} & \Lambda_u(m) & v_0 & \Lambda_\zeta |
| n_s | v_S | n_c | \log_m | r_1 + r_2 | \arg_\sigma |
|-----|------|-----|--------|-----------|------------|
| n_s | Z | 0 | 0 | 0 | 0 |
| n_c | Z | Z | Z | Z | 0 |
| r_1 + r_2 | \log_\sigma | R | R | 0 | \mathbb{Z} |
| r_2 | \arg_\sigma | R | R | 0 | \mathbb{Q} |
Step 5. We now compute the dual lattice by taking the inverse of the matrix (note that all the diagonal blocks are invertible). We obtain the following shape, where the red blocks a priori have coefficients in \( \mathbb{R} \). However, the coefficients in the \( \chi_m(g_i) \) columns (dual to the previous \( \log_m \) rows) represent a character on \( (\mathbb{Z}_F/m)^	imes \), which is a finite group, so they must be rationals with denominator divisible by the exponent of this group, and the coefficients in the \( k_\sigma \) columns (dual to the previous \( \arg_\sigma \) rows) represent characters on \( (\mathbb{R}/\mathbb{Z})^{r_2} \), so they must be integers.

\[
\begin{array}{c|cccc}
\text{n}_s & n_c & r_1 + r_2 & r_2 \\
\hline
\chi(p) & \chi_m(g_i) & \varphi_\sigma & k_\sigma \\
\text{n}_s & \text{Cl}_F & \mathbb{Q} & 0 & 0 \\
n_c & \text{Cl}(m) & \mathbb{Q} & \mathbb{Q} & 0 \\
r_1 + r_2 - 1 & (\hat{C}_m^1)_{k=0} & \mathbb{R} & \mathbb{Q} & \mathbb{R} \\
1 & v_0^\vee & \mathbb{R} & \mathbb{Q} & \mathbb{R} \\
r_2 & \hat{C}_m^1 & \mathbb{R} & \mathbb{Q} & \mathbb{R} \\
\end{array}
\]

Step 6. We remove \( v_0^\vee \), obtaining a basis of \( \Lambda^\vee \).

\[
\begin{array}{c|cccc}
\text{n}_s & n_c & r_1 + r_2 & r_2 \\
\hline
\chi(p) & \chi_m(g_i) & \varphi_\sigma & k_\sigma \\
n_s & \text{Cl}_F & \mathbb{Q} & 0 & 0 \\
n_c & \text{Cl}(m) & \mathbb{Q} & \mathbb{Q} & 0 \\
r_1 + r_2 - 1 & (\hat{C}_m^1)_{k=0} & \mathbb{R} & \mathbb{Q} & \mathbb{R} \\
r_2 & \hat{C}_m^1 & \mathbb{R} & \mathbb{Q} & \mathbb{R} \\
\end{array}
\]

A.2 Case with a CM subfield

In this section we assume that \( F \) contains a CM subfield. The implementation takes advantage of the following rationality result.

**Lemma 50.** Let \( K \) be a CM field and let \( w_K \) be the number of roots of unity in \( K \). Then for every \( u \in \mathbb{Z}_K^\times \) and \( \sigma : K \subseteq \mathbb{C} \), we have

\[
\frac{\arg(\sigma(u))}{2\pi} \in \frac{1}{2w_K}\mathbb{Z}.
\]

**Proof.** Let \( z = u/\bar{u} \in \mathbb{Z}_K^\times \). Then for every complex embedding \( \sigma \), we have \( |\sigma(z)| = 1 \). So \( z \) is a root of unity: \( z^{w_K} = 1 \). We obtain

\[
2\arg(\sigma(u)) = \arg(z) \in \frac{2\pi}{w_K}\mathbb{Z},
\]

hence the result. \(\square\)
For the remainder of this section, we assume that we are given \( K \) the maximal CM subfield of \( F \). We will write \( \tau \) for the complex embeddings of \( K \) and \( \sigma \) for the complex embeddings of \( F \). For every complex embedding \( \tau \) of \( K \), we let \( N_\tau(x) = \prod_{\sigma \mid \tau} \sigma(x) \).

We start by applying Steps 1 to 3 as in the previous case, obtaining the same shape, but we change the order of the columns.

\[
\begin{array}{|c|c|c|c|c|}
\hline
n_s & n_c & r_2 & r_1 + r_2 - 1 \\
\Lambda_S & \Lambda_{m,u} & \Lambda_\zeta & \Lambda_u(m) \\
\hline
n_s & v_S & Z & 0 & 0 & 0 \\
n_c & \log_m & Z & Z & 0 & 0 \\
r_1 + r_2 & \log_\sigma & \mathbb{R} & \mathbb{R} & 0 & \mathbb{R} \\
r_2 & \arg_\sigma & \mathbb{R} & \mathbb{R} & \mathbb{Q} & \mathbb{R} \\
\hline
\end{array}
\]

We now focus on the archimedean block, where we will apply extra column operations to exhibit the subgroup of almost-algebraic characters.

**Step 4’**. We introduce extra rows, parametrised by the complex embeddings \( \tau \) of \( K \), with values \( \arg(N_\tau(\epsilon))/2\pi \) for \( \epsilon \in \Lambda_u(m) \). Those values are in \( \frac{1}{2\pi} \mathbb{Z} \) by Lemma 50. We also select a subset of \( r_2 - r_2(K) \) complex embeddings such that the corresponding coordinates on \( \mathbb{R}^{r_2} \) are linearly independent, as linear forms, of the ones corresponding to \( N_\tau \). In the following matrices, we label the corresponding row by \( \arg' \).

\[
\begin{array}{|c|c|c|}
\hline
r_2 & r_1 + r_2 - 1 \\
\Lambda_\zeta & \Lambda_u(m) \\
\hline
r_1 + r_2 & \log_\sigma & 0 & \mathbb{R} \\
r_2 - r_2(K) & \arg' & \mathbb{Q} & \mathbb{R} \\
r_2(K) & \arg_{N_\tau} & \mathbb{Q} & \mathbb{Q} \\
\hline
\end{array}
\]

**Step 5’**. We apply a column HNF on the red blocks: this computes the subgroup \( \Lambda(\arg) \) of elements of \( \Lambda_\zeta + \Lambda_u(m) \) that have trivial \( \arg N_\tau \), and a complement \( \Lambda_{\arg} \) of \( \Lambda(\arg) \) in \( \Lambda_\zeta + \Lambda_u(m) \). We get the following shape.

\[
\begin{array}{|c|c|c|}
\hline
r_2(K) & n - r_2(K) - 1 \\
\Lambda_{\arg} & \Lambda(\arg) \\
\hline
r_1 + r_2 & \log_\sigma & \mathbb{R} & \mathbb{R} \\
r_2 - r_2(K) & \arg' & \mathbb{R} & \mathbb{R} \\
r_2(K) & \arg_{N_\tau} & \mathbb{Q} & 0 \\
\hline
\end{array}
\]

**Step 6’**. We insert \( v_0 \) as before.
\[
\begin{array}{c|ccc}
  & r_2(K) & n - r_2(K) & 1 \\
  \Lambda_{\text{arg}} & \Lambda(\text{arg}) & v_0 \\
  r_1 + r_2 & \log_\sigma & \mathbb{R} & \mathbb{R} & \mathbb{Z} \\
  r_2 - r_2(K) & \text{arg}' & \mathbb{R} & \mathbb{R} & 0 \\
  r_2(K) & \text{arg}_{\tau} & \mathbb{Q} & 0 & 0 \\
\end{array}
\]

Note that the block corresponding to columns \(\Lambda(\text{arg})\) and \(v_0\) and rows \(\log_\sigma\) and \(\text{arg}'\) is square and invertible.

**Step 7’**. Now we compute, as before, the dual basis of \(\Lambda^\vee\) by taking the inverse matrix.

On the archimedean block, the inverse of the matrix is a priori of the following shape.

\[
\begin{array}{c|ccc}
  & r_1 + r_2 & r_2 - r_2(K) & r_2(K) \\
  \varphi_\sigma & k'_\sigma & k_{\tau} \\
  r_2(K) & \text{a.a.} & 0 & 0 & \mathbb{Q} \\
  n - r_2(K) - 1 & \mathbb{R} & \mathbb{R} & \mathbb{R} \\
  1 & \mathbb{R} & \mathbb{R} & \mathbb{R} \\
\end{array}
\]

**Step 8’**. We delete \(v_0^\vee\) and change coordinates again to recover the usual parameters \(k_\sigma\). As in Step 5 above, the red blocks are a priori real but they must actually be integers.

\[
\begin{array}{c|cc}
  & r_1 + r_2 & r_2 \\
  \varphi_\sigma & k_\sigma \\
  r_2(K) & \text{a.a.} & 0 & \mathbb{Z} \\
  n - r_2(K) - 1 & \mathbb{R} & \mathbb{Z} \\
\end{array}
\]

In the end, we obtain the following shape for the matrix of characters.

\[
\begin{array}{c|cc|ccc}
  & n_s & n_c & r_1 + r_2 & r_2 \\
  \chi(p) & \chi_m(g_i) & \varphi_\sigma & k_\sigma \\
  n_s & \text{Cl}_F & \mathbb{Q} & 0 & 0 & 0 \\
  n_c & \text{Cl}(m) & \mathbb{Q} & \mathbb{Q} & 0 & 0 \\
  r_2(K) & \text{Cl}(m) & \mathbb{R} & \mathbb{Q} & 0 & \mathbb{Z} \\
  n - r_2(K) - 1 & \hat{C}_m & \mathbb{R} & \mathbb{Q} & \mathbb{R} & \mathbb{Z} \\
\end{array}
\]

This matrix is accessible as \texttt{gcharinit(bnf,mod)} [1] in our implementation in Pari/GP 2.15 [27]. As explained in section 2.5.1 we can recover the group of algebraic Hecke characters from \(\hat{C}_m\).