Supplementary material for the paper

Disk-based $k$-mer counting on a PC

by

Sebastian Deorowicz, Agnieszka Debudaj-Grabysz, and Szymon Grabowski

Contents

1 KMC usage

2 API
   2.1 CKmerAPI class ................................................................. 4
   2.2 CKMCF ile class .............................................................. 4

3 Example of API usage

4 Database format
   4.1 The .kmc.pre file structure ............................................... 9
   4.2 The .kmc.suf file structure .............................................. 10

5 Experimental results
   5.1 Test platforms ..................................................................... 11
   5.2 Parameters of programs ....................................................... 11
   5.3 Results .............................................................................. 12

6 Automatic setting of parameters in KMC

7 Selected components of the KMC algorithm (codes not shown in the main part of the paper)
1 KMC usage

KMC program constructs a database of statistics for input set of FASTQ files. This database can then be used from other software: directly via KMC API (described in Section 2) or by reading a textual file containing a list of \( k \)-mers and their related counters. This textual file can be obtained for a database by KMC-dump program, that is presented in Section 3 as a sample application of our KMC API. Section 4 describes the database format in detail for those interested in the low-level access to the data. Section 5 contains additional experimental results and a description of the parameters of execution of the examined programs. Section 6 contains description of how the automatic setting of parameters of KMC works.

Below we describe in detail the parameters and options of the KMC command-line tool, in version 0.3.

The general syntax is:

\[
\text{kmc [options] <input\_file\_name> <output\_file\_name> <working\_directory> [<working\_directory2> ...]}\]

or:

\[
\text{kmc [options] <@input\_file\_names> <output\_file\_name> <working\_directory> [<working\_directory2> ...]}\]

where the parameters are:

- **input\_file\_name** — a single file in FASTQ format (gzipped or not),
- **@input\_file\_names** — a file name with list of input files in FASTQ format (gzipped or not),
- **output\_file\_name** — the output database file; if such a file exists, it will be overwritten.

The configuration options comprise:

- **-v** — verbose mode (shows all parameter settings); default: false,
- **-k<\text{len}>** — \( k \)-mer length, \( k \) from 1 to \( \text{MAX\_K} \); default: 25,
- **-m<\text{size}>** — max amount of RAM in GB (from 4 to 1024); default: 32,
- **-p<\text{par}>** — set level of distribution (from 3 to 5); default: 4,
- **-f[a/q]** — input in FASTA format (-fa) or FASTQ format (-fq); default: FASTQ,
- **-q[value]** — use Quake’s compatible counting with [value] representing lowest quality; default: 33,
- **-ci<value>** — exclude \( k \)-mers appearing less than <value> times; default: 2,
- **-cs<value>** — maximal value of a counter; default: 255,
- **-cx<value>** — exclude \( k \)-mers appearing more of than <value> times; default: 1e9,
- **-b** — turn off transformation of \( k \)-mers into canonical form,
- **-sc<value>** — number of compacting threads (if not specified, it is set automatically based on the total number of computation threads, see the previous parameter),
- **-sf<value>** — number of FASTQ reading threads,
- **-sp<value>** — number of splitting threads,
- **-so<value>** — number of sorter threads,
• `-sr<value>` — number of threads per single sorter.

The parameters `-sc<value>`, `-sf<value>`, `-sp<value>`, `-so<value>`, and `-sr<value>` concern the internal work of KMC, i.e., their settings may affect the program’s processing speed, but won’t change its output. Not setting at least one parameter from this group makes KMC ignore them all.

Here are some usage examples.
```
kmc -k27 -m24 NA19238.fastq NA.res \data\kmc_tmp_dir\n
kmc -k27 -q -m24 @files.lst NA.res \data\kmc_tmp_dir\n```
2 API

In this section we describe two classes, CKmerAPI and CKMCFile. They can be used to gain access to the databases produced by KMC program.

2.1 CKmerAPI class

This class represents a *k*-mer. Its key methods are:

- CKmerAPI(uint32 length = 0) — constructor, that creates the array kmer_data of appropriate size,
- CKmerAPI(const CKmerAPI &kmer) — copy constructor,
- char get_symbol(unsigned int pos) — returns *k*-mer’s symbol at a given position (0-based),
- std::string to_string() — converts *k*-mer to string, using the alphabet ACGT,
- void to_string(char *str) — converts *k*-mer to string, using the alphabet ACGT; the function assumes that enough memory was allocated,
- void to_string(str::string &str) — converts *k*-mer to string, using the alphabet ACGT,
- bool from_string(std::string &str) — converts string (from alphabet ACGT) to *k*-mer,
- CKmerAPI() — destructor, releases the content of kmer_data array,
- overloaded operators: =, ==, <.

2.2 CKMCFile class

This class handles a *k*-mer database. Its key methods are:

- CKMCFile() — constructor,
- bool OpenForRA(std::string file_name) — opens two files: file_name with added extension “.kmc_pre” and “.kmc_suf”, reads their whole content to enable random access (in memory), and then closes them,
- bool OpenForListing(std::string file_name) — opens the file file_name with added extension “.kmc_pre” and allows to read the *k*-mers one by one (whole database is not loaded into memory),
- bool ReadNextKmer(CKmerAPI &kmer, float &count) — reads next *k*-mer to kmer and updates its count; the return value is bool; true as long as not eof-of-file (available only when database is opened in listing mode),
- bool Close() — if the file was opened for random access, the allocated memory for its content is released; if the file was opened for listing, the allocated memory for its content is released and the “.kmer” file is closed,
- bool SetMinCount(uint32 x) — set the minimum counter value for *k*-mers; if a *k*-mer has count below x, it is treated as non-existent,
- uint32 GetMinCount(void) — returns the value (uint32) set with SetMinCount,
• `bool SetMaxCount(uint32 x)` — set the maximum counter value for $k$-mers; if a $k$-mer has count above $x$, it is treated as non-existent,

• `uint32 GetMaxCount(void)` — returns the value (uint32) set with `SetMaxCount`,

• `uint64 KmerCount(void)` — returns the number of $k$-mers in the database (available only for databases opened in random access mode),

• `uint32 KmerLength(void)` — returns the $k$-mer length in the database (available only for databases opened in random access mode),

• `bool RestartListing(void)` — sets the cursor for listing $k$-mers from the beginning of the file (available only for databases opened in listing mode). The method `OpenForListing(std::string file_name)` invokes it automatically, but it can be also called by a user,

• `bool Eof(void)` — returns true if all $k$-mers have been listed,

• `bool CheckKmer(CKmerAPI &kmer, float &count)` — returns true if kmer exists in the database and set its count if the answer is positive (available only for databases opened in random access mode),

• `bool IsKmer(CKmerAPI &kmer)` — returns true if kmer exists (available only for databases opened in random access mode),

• `void ResetMinMaxCounts(void)` — sets min_count and max_count to the values read from the database,

• `bool Info(uint32 &kmer_length, uint32 &mode, uint32 &counter_size, uint32 &lut_prefix_length, uint32 &min_count, uint32 &max_count, uint64 &total_kmers)` — gets current parameters from the $k$-mer database,

• `CKMCFile()` — destructor.
3 Example of API usage

The KMC-dump application (Figs. S1 and S2) shows how to list and print $k$-mers with at least $min\_count$ and at most $max\_count$ occurrences in the database. Fig. S1 presents parsing the command-line parameters, including -ci<$value>$ and -cx<$value>$. Input and output file names are also expected. The code in Fig. S2 is for actual database handling. This database is represented by a CKMCFile object, which opens an input file for $k$-mer listing (the method bool OpenForListing(std::string file_name) is invoked). The parameter of the method SetMinCount (SetMaxCount) must be not smaller (not greater) than the corresponding parameter -ci (-cx) with which KMC was invoked (otherwise, nothing will be listed). The listed $k$-mers are in the form like: AAACACCGT\t$<$value>\$

where the first part is the $k$-mer in natural representation, which is followed by a tab character, and its associated value (integer or float). (Such format is compatible with Quake, a widely used tool for sequencing error correction.) Note that, if needed, one can easily modify the output format, changing the lines 39 and 41 in Fig. S2.
```cpp
#include <iostream>
#include "./kmc_api/kmc_file.h"

void print_info(void);

int _main(int argc, char* argv[])
{
    CKMCFile kmer_database;
    int i;
    uint32 min_count_to_set = 0;
    uint32 max_count_to_set = 0;
    std::string input_file_name;
    std::string output_file_name;

    FILE * out_file;

    // Parse input parameters
    //---
    if (argc < 3)
    {
        print_info();
        return EXIT_FAILURE;
    }
    for (i = 1; i < argc; ++i)
    {
        if (argv[i][0] == '-')
        {
            if (strncmp(argv[i], "-ci", 3) == 0)
                min_count_to_set = atoi(&argv[i][3]);
            else if (strncmp(argv[i], "-cx", 3) == 0)
                max_count_to_set = atoi(&argv[i][3]);
        } else
            break;
    }
    if (argc - i < 2)
    {
        print_info();
        return EXIT_FAILURE;
    }
    input_file_name = std::string(argv[i++]);
    output_file_name = std::string(argv[i]);
    if ((out_file = fopen (output_file_name.c_str(), "w")) == NULL)
    {
        print_info();
        return EXIT_FAILURE;
    }
    setvbuf(out_file, NULL, _IOFBF, 1 << 24);
    ...
```

Figure S1: First part of KMC-dump application
// Open kmer database for listing and print kmers within min_count and max_count

if (!kmer_database.OpenForListing(input_file_name))
{
    print_info();
    return EXIT_FAILURE;
}
else
{
    uint32 _kmer_length;
    uint32 _mode;
    uint32 _counter_size;
    uint32 _lut_prefix_length;
    uint32 _min_count;
    uint32 _max_count;
    uint64 _total_kmers;

    kmer_database.info(_kmer_length, _mode, _counter_size, _lut_prefix_length, _min_count, _max_count, _total_kmers);

    float counter;
    std::string str;

    CKmerAPI kmer_object(_kmer_length);

    if (min_count_to_set)
        if (!kmer_database.SetMinCount(min_count_to_set))
            return EXIT_FAILURE;
    if (max_count_to_set)
        if (!kmer_database.SetMaxCount(max_count_to_set))
            return EXIT_FAILURE;

    while (kmer_database.ReadNextKmer(kmer_object, counter))
    {
        kmer_object.to_string(str);
        if (_mode)
            fprintf(out_file, "%s\t%f\n", str.c_str(), counter);
        else
            fprintf(out_file, "%s\t%d\n", str.c_str(), (int)counter);
    }

    fclose(out_file);
}

return EXIT_SUCCESS;

// Print execution options

void print_info(void)
{
    std::cout << "KMC dump ver. " << KMC_VER << "( " << KMC_DATE << " )" << std::endl;
    std::cout << "Usage: nkmc_dump [ options ] <kmc_database> <output_file>" << std::endl;
    std::cout << "Parameters:" << std::endl;
    std::cout << "-ci<value> - print k-mers appearing less than <value> times" << std::endl;
    std::cout << "-cx<value> - print k-mers appearing more of than <value> times" << std::endl;
};

Figure S2: Second part of KMC-dump application
4 Database format

The KMC application creates output files with two extensions:

- `.kmc.pre` — information on $k$-mer prefixes (plus some other data) are stored here,
- `.kmc.suf` — information on $k$-mer suffixes and the related counters are stored here.

All integers in the KMC output files are stored in LSB (least significant byte first) format.

4.1 The `.kmc.pre` file structure

A `.kmc.pre` file contains, in the order:

- [marker],
- [data],
- [header],
- [header position],
- [marker] (another copy, to signal the file is not truncated).

[marker]
4 bytes with the letters: KMCP.

[header position]
An integer consisting of the last 4 bytes in the file. Its contains the relative position of the beginning of the field [header]. After opening the file, one should do the following:

1. Read the first 4 bytes and check if they contain the letters KMCP.
2. Jump to position 4 bytes back from end of file and read the header position $x$.
3. Jump to position $x + 4$ bytes back from end of file and read the header.
4. Read [data].

[header]
The header contains fields describing the file `.kmc.pre`:

- `uint32 kmer_length` — $k$-mer length,
- `uint32 mode` — mode: 0 (occurrence count) or 1 (counting according to Quake quality),
- `uint32 counter_size` — counter field size: for mode 0 it is 1, 2, 3, or 4; for mode 1 it is always 4,
- `uint32 lut_prefix_length` — the length (in symbols) of the prefix cut off from $k$-mers; it is invariant of the scheme that 4 divides $kmer_length - lut_prefix_length$,
• `uint32 min_count` — minimum number of k-mer occurrences to write in the database (if the counter is smaller, the k-mer data are not written),

• `uint32 max_count` — maximum number of k-mer occurrences to write in the database (if the counter is greater, the k-mer data are not written),

• `uint64 total_kmers` — total number of k-mers in the database,

• `uint32 tmp[8]` — not used in the current version.

Here are k-mer prefix data. More precisely, it is an array of `uint64` elements, of size \(4\text{lut}_{\text{prefix}}\text{.length}\). Position \(x\) in the array stores the number of different k-mers whose prefix of length \(\text{lut}_{\text{prefix}}\text{.length}\) is (in binary) less than \(x\). DNA symbols are encoded as follows: \(A \leftarrow 0, C \leftarrow 1, G \leftarrow 2, T \leftarrow 3\). For example, if the queried k-mer is `ACGACTGAT` and \(\text{lut}_{\text{prefix}}\text{.length} = 4\), then we cut off the first 4 symbols, i.e., the prefix `ACGA`, which is interpreted in binary as \(x = 24\) (since \(0 \times 2^6 + 1 \times 2^4 + 2 \times 2^2 + 0 \times 2^0 = 24\)). Now we look into “data” at locations \(x\) and \(x + 1\) to read, e.g., 1523 and 1685. This means that in the file .kmc.suf in records from 1523 to 1685 – 1 there are suffixes of the k-mers with prefix ACGA. Having got this range, we can now binary search the suffix CTGAT.

4.2 The .kmc.suf file structure

A .kmc.suf file contains, in order:

• [marker],

• [data],

• [marker] (another copy, to signal the file is not truncated).

The k-mers are stored with their leftmost symbol first, packed into bytes. For example, CCACAAAT is represented as 0x51 (for CCAC), 0x03 (for AAAT). Integers are stored according to the LSB (little endian) convention, floats are stored in the same way as they are stored in the memory.

[marker]

4 bytes with the letters: KMCS.

[data]

An array `record_t records[total_kmers].`

`total_kmers` is a value taken from the .kmc_pre file.

`record_t` is a type representing a k-mer. Its first field is the k-mer suffix string, stored on \((kmer_{length} – \text{lut}_{\text{prefix}}\text{.length})/4\) bytes. The next field is `counter_size`, with the number of bytes used by the counter, which is either a 1...4-byte integer, or a 4-byte float.
5 Experimental results

5.1 Test platforms

K-mer Counter (KMC), was implemented in C++11, using gcc compiler (version 4.7.1) for the linux build and Microsoft Visual Studio 2012 for the Windows build. Experiments comprise also the tools Jellyfish and BFCounter.

Two test machines were used. One was a 4 AMD Opteron™ 6136 2.4 GHz CPUs (32 cores) server with 128 GB RAM, and fast RAID-0 disk matrix of total size 2.0 TB. The other was a “home” PC, with 6-core AMD Phenom II 1090 3.2 GHz CPU, 16 GB RAM and 3 SATA HDD of sizes 2 TB each. The hard disks at the PC machine were: two Seagate Barracuda Green ST2000DL003 with 5,900 rpm and one Hitachi Deskstar 7K3000 with 7,200 rpm.

5.2 Parameters of programs

BFCounter

The main problem with running BFCounter (ver. 0.2) concerns proper estimation of the number of distinct \( k \)-mers. In the experiments we set this parameter to be about 20 percent greater than the actual number of distinct \( k \)-mers (which was known from a KMC test).

Jellyfish

Jellyfish (ver. 0.5) requires to give as a parameter the size of the hash table, which should be a bit larger than the total number of distinct \( k \)-mers. Similarly as for BFCounter we passed the value about 20 percent greater than the number of distinct \( k \)-mers (known from prior KMC execution). Unfortunately, for some cases (denoted by asterisk in the tables with experimental results), it resulted in Jellyfish crash due to the insufficient amount of RAM (128 GB of RAM was available in the test machine). In such cases we halved the parameter value and ran Jellyfish again, which produced two output files. These files should be merged then, but unfortunately, Jellyfish was unable to do that due to insufficient amount of RAM.

To allow right comparison with BFCounter (which does not count unique \( k \)-mers), we set in the execution of Jellyfish that it should not count \( k \)-mers with count value less than 2.0.

Probably due to some bug, when the output file name specified in the command line is shorter than 10 letters, Jellyfish often dramatically slows down (2–3 times). Thus, we always used long file names.

DSK

DSK (ver. 1.4993) was executed with default parameters that means 5 GB limit of RAM. The \( k \)-mers occurring less than 2 times were excluded.

KMC

KMC was executed with setting that \( k \)-mers occurring less than 2 times should not be counted. The parameter \( -p \) was set to 3 for two smallest data sets, to 4 for *Caenorhabditis elegans* data set and to 5 for two human data sets.
5.3 Results

The results presented at the following pages are for 5 data sets characterized in the main part of the paper. For 3 data sets examined in the main paper we duplicate here the tables and provide extra results (for other values of $k$).
Table S1: $k$-mers counting results for *D. ananassae* (8.7 GB FASTQ file or 6 gzipped FASTQ files of total size 1.8 GB). RAM and disk spaces are in GB (1GB=$2^{30}$B). Time is in seconds. The programs were used for the number of threads adjusted to the number of cores to achieve maximum speed.

| Algorithm | $k = 22$ | $k = 25$ | $k = 28$ | $k = 31$ | $k = 40$ | $k = 55$ | $k = 70$ |
|-----------|----------|----------|----------|----------|----------|----------|----------|
|           | Space    | Time     | Space    | Time     | Space    | Time     | Space    | Time     |
| Classic counters |          |          |          |          |          |          |          |          |
| BFCounter  | 0.6/ 0   | 2,157    | 0.6/ 0   | 2,082    | failed   | failed   | failed   | failed   |
| Jellyfish  | 10/ 0    | 42       | 10/ 0    | 48       | 11/ 0    | 44       | 11/ 0    | 44       |
| KMC        | 32/ 4    | 43       | 32/ 5    | 44       | 32/ 7    | 48       | 32/ 7    | 49       |
| KMC        | 16/ 4    | 45       | 16/ 5    | 44       | 16/ 7    | 47       | 16/ 7    | 44       |
| KMC$_{gz}$ | 32/ 4    | 47       | 32/ 5    | 45       | 32/ 7    | 49       | 32/ 7    | 50       |
| KMC$_{gz}$ | 16/ 4    | 46       | 16/ 5    | 49       | 16/ 7    | 49       | 16/ 7    | 48       |
| 6-core PC  |          |          |          |          |          |          |          |          |
| KMC        | 11/ 4    | 103      | 11/ 5    | 130      | 11/ 6    | 135      | 11/ 7    | 143      |
| KMC$_{gz}$ | 11/ 4    | 79       | 11/ 5    | 85       | 11/ 6    | 89       | 11/ 7    | 101      |

Quake-compatible counters

| Algorithm | $k = 22$ | $k = 25$ | $k = 28$ | $k = 31$ | $k = 40$ | $k = 55$ | $k = 70$ |
|-----------|----------|----------|----------|----------|----------|----------|----------|
|           | Space    | Time     | Space    | Time     | Space    | Time     | Space    | Time     |
| BFCounter  | 0.7/ 0   | 2,126    | 0.7/ 0   | 2,049    | 0.7/ 0   | 1,889    | 0.7/ 0   | 1,850    |
| Jellyfish  | 10/ 0    | 54       | 10/ 0    | 52       | 10/ 0    | 53       | 10/ 0    | 52       |
| KMC        | 32/12    | 77       | 32/12    | 78       | 32/13    | 91       | 32/14    | 82       |
| KMC        | 16/12    | 61       | 16/12    | 61       | 16/13    | 62       | 16/14    | 67       |
| KMC$_{gz}$ | 32/12    | 85       | 32/12    | 85       | 32/13    | 86       | 32/14    | 86       |
| KMC$_{gz}$ | 16/12    | 94       | 16/12    | 65       | 16/13    | 67       | 16/14    | 70       |
| 6-core PC  |          |          |          |          |          |          |          |          |
| KMC        | 11/12    | 171      | 11/12    | 187      | 11/13    | 190      | 11/13    | 194      |
| KMC$_{gz}$ | 11/12    | 173      | 11/12    | 177      | 11/13    | 171      | 11/13    | 178      |
### Table S2: k-mers counting results for *Caenorhabditis elegans* (16.4 GB FASTQ file or 2 gzipped FASTQ files of total size 4.6 GB). Test methodology and column description are just as for Table S1.

| Algorithm  | k = 22 | k = 25 | k = 28 | k = 31 | k = 40 | k = 55 | k = 70 |
|------------|--------|--------|--------|--------|--------|--------|--------|
|            | Space  | Time   | Space  | Time   | Space  | Time   | Space  | Time   |
| BFCounter  | 4/ 0   | 10,407 | 4/ 0   | 10,142 | failed | failed | failed | failed |
| Jellyfish  | 21/ 0  | 88     | 22/ 0  | 130    | 22/ 0  | 89     | 23/ 0  | 89     |
| KMC        | 32/10  | 105    | 32/14  | 99     | 32/18  | 115    | 32/20  | 123    |
| KMC        | 16/11  | 93     | 16/14  | 84     | 16/18  | 93     | 16/21  | 105    |
| KMC        | 32/10  | 134    | 32/14  | 126    | 32/18  | 138    | 32/20  | 146    |
| KMC        | 16/11  | 119    | 16/14  | 127    | 16/18  | 121    | 16/21  | 126    |
| KMC        | 11/11  | 274    | 11/14  | 316    | 11/18  | 343    | 11/21  | 379    |
| KMC        | 11/11  | 233    | 11/14  | 267    | 11/18  | 333    | 11/21  | 359    |
| KMC        | 11/32  | 562    | 11/34  | 581    | 11/37  | 635    | 11/40  | 678    |
| KMC        | 11/32  | 555    | 11/34  | 581    | 11/37  | 627    | 11/40  | 673    |

#### Classic counters

| 32-core server | 6-core PC |
|----------------|-----------|
| BFCounter      | 4/ 0      | 10,349   |
| Jellyfish      | 24/ 0     | 143      |
| KMC            | 32/31     | 173      |
| KMC            | 16/32     | 166      |
| KMC            | 32/31     | 184      |
| KMC            | 16/32     | 168      |
| KMC            | 11/32     | 562      |
| KMC            | 11/32     | 555      |

#### Quake-compatible counters

| 32-core server | 6-core PC |
|----------------|-----------|
| BFCounter      | 4/ 0      | 10,044   |
| Jellyfish      | 25/ 0     | 154      |
| KMC            | 32/37     | 173      |
| KMC            | 16/37     | 179      |
| KMC            | 32/37     | 186      |
| KMC            | 16/37     | 173      |
| KMC            | 11/37     | 581      |
| KMC            | 11/37     | 581      |

#### 22-core counter in 16-core PC

| 32-core server | 6-core PC |
|----------------|-----------|
| BFCounter      | 4/ 0      | 10,044   |
| Jellyfish      | 24/ 0     | 143      |
| KMC            | 32/37     | 173      |
| KMC            | 16/37     | 179      |
| KMC            | 32/37     | 186      |
| KMC            | 16/37     | 173      |
| KMC            | 11/37     | 581      |
| KMC            | 11/37     | 581      |
Table S3: k-mers counting results for Zea mays (45.9 GB FASTQ file or 108 gzipped FASTQ files of total size 16.3 GB). Test methodology and column description are just as for Table S1.

| Algorithm | k = 22 | k = 25 | k = 28 | k = 31 | k = 40 | k = 55 | k = 70 |
|-----------|--------|--------|--------|--------|--------|--------|--------|
|           | Space  | Time   | Space  | Time   | Space  | Time   | Space  | Time   |
|           |        |        |        |        |        |        |        |        |
| Classic counters |        |        |        |        |        |        |        |        |
| BFCounter | 12/0   | 38,050 | 13/0   | 37,509 | failed | failed | failed | failed |
| Jellyfish | 60/0   | 358    | 63/0   | 524    | 66/0   | 378    | 70/0   | 395    |
| KMC       | 32/40  | 328    | 32/56  | 346    | 32/72  | 433    | 32/82  | 471    | 32/125 | 792   | 32/183 | 1,082 | 32/230 | 1,576 |
| KMC       | 16/42  | 287    | 16/56  | 302    | 16/74  | 385    | 16/87  | 455    | 16/125 | 785   | 16/189 | 1,107 | 16/235 | 1,608 |
| KMC$^\text{gz}$ | 32/40 | 423    | 32/56  | 365    | 32/72  | 480    | 32/82  | 532    | 32/125 | 804   | 32/183 | 1,060 | 32/230 | 1,530 |
| KMC$^\text{gz}$ | 16/41 | 338    | 16/56  | 310    | 16/74  | 347    | 16/86  | 424    | 16/125 | 736   | 16/188 | 1,037 | 16/235 | 1,566 |
| 6-core PC |
| KMC       | 11/42  | 809    | 11/58  | 893    | 11/75  | 1,138  | 11/89  | 1,324  | 11/126 | 2,151 | 11/188 | 3,081 | 11/238 | 4,846 |
| KMC$^\text{gz}$ | 11/42 | 734    | 11/58  | 889    | 11/75  | 1,180  | 11/89  | 1,405  | 11/126 | 2,160 | 11/188 | 3,227 | 11/238 | 4,846 |
| Quake-compatible counters |        |        |        |        |        |        |        |        |
| BFCounter | 34/0   | 44,025 | 40/0   | 44,838 | failed | failed | failed | failed |
| Jellyfish | 72/0   | 727    | 75/0   | 662    | 78/0   | 726    | 82/0   | 741    |
| KMC       | 32/117 | 690    | 32/133 | 750    | 32/150 | 824    | 32/160 | 888    | 32/198 | 1,189 | 32/252 | 1,479 | 32/298 | 2,048 |
| KMC       | 16/121 | 681    | 16/135 | 757    | 16/150 | 782    | 16/165 | 902    | 16/198 | 1,213 | 16/257 | 1,530 | 16/303 | 2,166 |
| KMC$^\text{gz}$ | 32/117 | 675    | 32/133 | 711    | 32/150 | 787    | 32/160 | 850    | 32/198 | 1,142 | 32/252 | 1,422 | 32/297 | 1,989 |
| KMC$^\text{gz}$ | 16/121 | 666    | 16/135 | 698    | 16/150 | 801    | 16/164 | 837    | 16/198 | 1,144 | 16/257 | 1,479 | 16/303 | 2,222 |
| 6-core PC |
| KMC       | 11/123 | 1,866  | 11/134 | 2,125  | 11/150 | 2,370  | 11/167 | 2,532  | 11/198 | 3,360 | 11/256 | 4,574 | 11/304 | 6,327 |
| KMC$^\text{gz}$ | 11/123 | 1,951  | 11/135 | 2,203  | 11/150 | 2,346  | 11/167 | 2,658  | 11/198 | 3,611 | 11/256 | 4,558 | 11/304 | 6,496 |
Table S4: $k$-mers counting results for *Homo sapiens* NA19238 individual (353 GB FASTQ file or 463 gzipped FASTQ files of total size 116.6 GB). Test methodology and column description are just as for Table S1. The asterisk signs (for Jellyfish) denote that two separate databases were constructed by Jellyfish due to the memory limit of the machine (128 GB RAM) and Jellyfish reported that to merge these databases it needs more RAM, so these times are underestimated.

| Algorithm | $k = 22$ | $k = 25$ | $k = 28$ | $k = 31$ |
|-----------|----------|----------|----------|----------|
|           | Space    | Time     | Space    | Time     | Space    | Time     |
| **Classic counters** |          |          |          |          |          |          |
| 32-core server |          |          |          |          |          |          |
| BFCounter  | 46/0     | 114,083  | 41/0     | 99,468   | failed   | failed   |
| Jellyfish  | 50/0     | 2,303    | 64/0     | 2,258    | 75/0     | 2,208    |
| KMC        | 32/104   | 1,405    | 32/130   | 1,488    | 32/133   | 1,522    |
|            | 16/107   | 1,548    | 16/131   | 1,657    | 16/141   | 1,684    |
| KMC$^{**}$ | 32/104   | 1,040    | 32/129   | 1,066    | 32/132   | 1,035    |
| KMC$^{**}$ | 16/107   | 1,278    | 16/130   | 1,631    | 16/141   | 1,662    |
| 6-core PC  |          |          |          |          |          |          |
| KMC        | 11/107   | 3,482    | 11/128   | 3,442    | 11/138   | 3,584    |
| KMC$^{**}$ | 11/107   | 2,198    | 11/128   | 2,206    | 11/138   | 2,303    |
| **Quake-compatible counters** |          |          |          |          |          |          |
| 32-core server |          |          |          |          |          |          |
| BFCounter  | 70/0     | 171,888  | 72/0     | 180,861  | failed   | failed   |
| Jellyfish  | 100/0    | 4,339    | 57/230   | 2,891$^*$| 64/192   | 3,246$^*$|
| KMC        | 32/311   | 2,585    | 32/302   | 2,467    | 32/282   | 2,347    |
| KMC$^{**}$ | 16/315   | 2,615    | 16/305   | 2,730    | 16/283   | 2,592    |
| KMC$^{**}$ | 32/310   | 2,071    | 32/302   | 1,995    | 32/282   | 1,880    |
| KMC$^{**}$ | 16/318   | 2,273    | 16/304   | 2,611    | 16/283   | 2,015    |
| 6-core PC  |          |          |          |          |          |          |
| KMC        | 11/316   | 5,538    | 11/298   | 5,533    | 11/277   | 5,184    |
| KMC$^{**}$ | 11/316   | 5,570    | 11/298   | 5,060    | 11/277   | 4,708    |
Table S5: \( k \)-mers counting results for \( \textit{Homo sapiens} \) HG02057 individual (208 GB FASTQ file or 6 gzipped FASTQ files of total size 659 GB). Test methodology and column description are just as for Table S1. The asterisk signs (for Jellyfish) denote that two separate databases were constructed by Jellyfish due to the memory limit of the machine (128 GB RAM) and Jellyfish reported that to merge these databases it needs more RAM, so these times are underestimated.

| \( k \) | Algorithm | Space | Time |
|-------|----------|-------|------|
| 22    | Jellyfish | 27/0  | 1,175 |
|       | KMC      | 32/38 | 11/234 3,024 |
|       | KMC\(^\d\) | 16/134 | 11/234 3,024 |
| 25    | Jellyfish | 32/0  | 1,375 |
|       | KMC      | 32/38 | 12/234 3,024 |
|       | KMC\(^\d\) | 16/134 | 12/234 3,024 |
| 28    | Jellyfish | 32/0  | 1,375 |
|       | KMC      | 32/38 | 12/234 3,024 |
|       | KMC\(^\d\) | 16/134 | 12/234 3,024 |
| 31    | Jellyfish | 32/0  | 1,375 |
|       | KMC      | 32/38 | 12/234 3,024 |
|       | KMC\(^\d\) | 16/134 | 12/234 3,024 |
| 40    | Jellyfish | 32/0  | 1,375 |
|       | KMC      | 32/38 | 12/234 3,024 |
|       | KMC\(^\d\) | 16/134 | 12/234 3,024 |
| 55    | Jellyfish | 32/0  | 1,375 |
|       | KMC      | 32/38 | 12/234 3,024 |
|       | KMC\(^\d\) | 16/134 | 12/234 3,024 |
| 70    | Jellyfish | 32/0  | 1,375 |
|       | KMC      | 32/38 | 12/234 3,024 |
|       | KMC\(^\d\) | 16/134 | 12/234 3,024 |

Note: Times are underestimated due to the memory limit of the machine (128 GB RAM) and Jellyfish reported that to merge these databases it needs more RAM.
6 Automatic setting of parameters in KMC

The automatic setting of parameters mechanism tries to allocate the available resources (i.e., CPU cores) in the best possible way. The optimal number of threads for the parts of the algorithm is, however, hard to obtain, since it depends on many things, like the compression method of input files, the number and speed of disks, etc. Thus, our automatic mechanism is obviously suboptimal, nevertheless, experiments show that it performs reasonably well. If the results are unsatisfactory, the KMC user can specify these parameters from command line.

The most important factor of the mechanism is the number of available cores (possibly overridden if the user specifies it with -s? parameters).

Table S6: Automatic selection of internal parameters in KMC. $T$ is the number of cores, $D$ the number of working directories, $F$ the number of input files

| FASTQ readers | Splitters | Compactors | Disk writers and Bin writers | Sorters | Sorting threads per sorter |
|---------------|-----------|------------|------------------------------|---------|---------------------------|
| $N_F$         | $N_{Sp}$  | $N_C$      | $N_D$                        | $N_{So}$| $N_{St}$                  |

| Small PC: $T < 6$ |
|-------------------|
| $\min(F, T)$      | $\max(1, T - 1)$ | $\max(1, T - 1)$ | $D$ | 1 | $T$ |

| Large PC: $6 \leq T \leq 12$ |
|-----------------------------|
| $\min(F, T)$                | $\max(1, T - 1)$ | $\max(1, T - 1)$ | $D$ | 2 | $T/2 + 1$ |

| Small server: $12 < T \leq 24$ |
|---------------------------------|
| $\min(F, T/3)$                 | $(T - N_F) / 2$ | $T - N_F - N_{Sp} - 1$ | $D$ | $T/6$ | $T/N_{So}$ |

| Large server: $24 < T$ |
|------------------------|
| $\min(F, T/3)$         | $(T - N_F) / 2$ | $T - N_F - N_{Sp} - 1$ | $D$ | $T/8$ | $T/N_{So}$ |
7 Selected components of the KMC algorithm (codes not shown in the main part of the paper)

**FASTQ-READER**(*FL_queue*)

Input: *FL_queue* — queue with list of input file names

| Line | Code                                                                 | Description |
|------|----------------------------------------------------------------------|-------------|
| 1    | while (file_name ← FL_QUEUE.pop()) not empty do                      |             |
| 2    | file ← Open file file_name                                           |             |
| 3    | while (block ← file.read_next_block()) not empty do                  |             |
| 4    | FASTQ_parts_queue.push(block)                                       |             |
| 5    | Close file                                                           |             |

Figure S3: Reading FASTQ/FASTA files. Each block, of several megabytes, is rounded to the nearest complete read. The blocks are inserted into a FASTQ parts queue.

**COMPACTOR(...)**

Input: ... — queue with parts of FASTQ/FASTA files

| Line | Code                                                                 | Description |
|------|----------------------------------------------------------------------|-------------|
| 1    | while (bin_id, bin_package ← Bin_part_packages_queue.pop()) not empty and |             |
|      | not Bin_part_packages_queue.finished() do                            |             |
| 2    | if bin_package not empty then                                         |             |
| 3    | Try to compact bin_package on \( p_2 + 4 \) symbols                  |             |
| 4    | Put (possibly compacted) bin_package to Compacted_buffers_queue       |             |

Figure S4: Algorithm of extra compacting bin packages prior to writing to disk. The output of this stage are compacted packages.

**DISK-WRITER(...)**

Input: ... — queue with parts of FASTQ/FASTA files

| Line | Code                                                                 | Description |
|------|----------------------------------------------------------------------|-------------|
| 1    | while (bin_id, buffer ← Compacted_buffer_queue.pop()) not empty and   |             |
|      | not Compacted_buffer_queue.finished() do                             |             |
| 2    | if buffer not empty then                                             |             |
| 3    | Write buffer to disk file related with bin_id                         |             |

Figure S5: Disk-writer algorithm
**Bin reader(...)**

Input: ... — queue with parts of FASTQ/FASTA files

1. while (bin_id ← Bin_ids_queue.pop()) not empty do
2.   bin_data ← read whole file related to bin_id
3.   Bins_queue.push(bin_id, bin_data)

Figure S6: Bin reader algorithm

---

**Completer(...)**

Input: ... — queue with parts of FASTQ/FASTA files

1. while (bin_id, buffers ← Sorted_and_compacted_bins_priority_queue.pop()) not empty and
   not Sorted_and_compacted_bins_priority_queue.finished() do
2.   if buffers not empty then
3.     Store buffers to two output files

Figure S7: Completer algorithm