Second-Order Asymptotics for the Gaussian MAC with Degraded Message Sets

Jonathan Scarlett and Vincent Y. F. Tan

Abstract

This paper studies the second-order asymptotics of the Gaussian multiple-access channel with degraded message sets. For a fixed average error probability $\varepsilon \in (0,1)$ and an arbitrary point on the boundary of the capacity region, we characterize the speed of convergence of rate pairs that converge to that boundary point for codes that have asymptotic error probability no larger than $\varepsilon$. We do so by elucidating clear relationships between the rate pairs achievable at large blocklengths and the local second-order asymptotics, i.e. the second-order behavior of these rate pairs in the neighborhood of a boundary point. We provide a numerical example to illustrate how the angle of approach to a boundary point affects the second-order coding rate. This is the first conclusive characterization of the second-order asymptotics of a network information theory problem in which the capacity region is not a polygon.

Index Terms

Gaussian multiple-access channel, Degraded message sets, Superposition coding, Strong converse, Finite block-lengths, Second-order coding rates, Dispersion.

I. INTRODUCTION

In this paper, we revisit the Gaussian multiple-access channel (MAC) with degraded message sets (DMS). This is a communication model in which two independent messages are to be sent from two sources to a common destination. One encoder, the cognitive or informed encoder, has access to both messages, while the uninformed encoder only has access to its own message. Both transmitted signals are assumed to be power limited, and their sum is corrupted by additive white Gaussian noise (AWGN). See Fig. 1.

The capacity region, i.e. the set of all pairs of achievable rates, is well-known (e.g. see [1, Ex. 5.18(b)]), and is given by the set of rate pairs $(R_1, R_2)$ satisfying

\[ R_1 \leq C((1 - \rho^2)S_1) \]

\[ R_1 + R_2 \leq C(S_1 + S_2 + 2\rho \sqrt{S_1S_2}) \]

for some $\rho \in [0,1]$, where $S_1$ and $S_2$ are the admissible transmit powers, and $C(x) := \frac{1}{2}\log(1 + x)$ is the Gaussian capacity function. The capacity region for $S_1 = S_2 = 1$ is illustrated in Fig. 2. The boundary is parametrized by $\rho$, and the direct part is proved using superposition coding [2].

While the capacity region is well-known, there is substantial motivation to understand the second-order asymptotics for this problem. For any given point $(R_1^*, R_2^*)$ on the boundary of the capacity region, we study the rate of convergence to that point for an $\varepsilon$-reliable code. More precisely, we characterize the set of all $(L_1, L_2)$ pairs, known as second-order coding rates [3]–[6], for which there exist sequences of codes whose asymptotic error probability does not exceed $\varepsilon$, and whose code sizes $M_{1,n}$ and $M_{2,n}$ behave as

\[ \log M_{j,n} \geq nR_j^* + \sqrt{n}L_j + o(\sqrt{n}), \quad j = 1, 2. \]

This study allows us to understand the fundamental tradeoffs between the rates of transmission and average error probability from a perspective different from the study of error exponents. Here, instead of fixing a pair of rates and studying the exponential decay of the error probability $\varepsilon$, we fix $\varepsilon$ and study the speed at which a sequence of rate pairs approaches an information-theoretic limit as the blocklength grows.

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A. Related Work

The study of the second-order asymptotics of channel coding for discrete memoryless channels was initiated by Strassen [7]. For the single-user AWGN channel with (equal or maximal) power constraint $S$, a specialization of our model with $M_{2,n} = 1$, Hayashi [4] and Polyanskiy et al. [8] showed that the optimum (highest) second-order coding rate is $\sqrt{V(S)\Phi^{-1}(\varepsilon)}$, where $V(x) := \frac{x(x+2)}{2(x+1)^2}$ is the Gaussian dispersion function. More precisely, Polyanskiy et al. [8, Thm. 54] showed the asymptotic expansion

$$\log M^*(n,\varepsilon) = nC(S) + \sqrt{nV(S)\Phi^{-1}(\varepsilon)} + g_n,$$

where $M^*(n,\varepsilon)$ is the maximum size of a length-$n$ block code with (either average or maximum) error probability $\varepsilon$, and $O(1) \leq g_n \leq \frac{1}{2}\log n + O(1)$. In fact, the expression for $V(S)$ was already known to Shannon [9, Sec. X], who analyzed the reliability function of the AWGN channel for rates close to capacity.

There have been numerous attempts to study the finite blocklength behavior and second-order asymptotics for MACs [10]–[18], but most of these works focus on inner bounds (the direct part). The development of tight and easily-evaluated converse bounds remains more modest, and those available do not match the direct part in general or are very restrictive (e.g. product channels were considered in [18]). We will see that the assumption of Gaussianity of the channel model together with the degradedness of the message sets allows us to circumvent some of the difficulties in proving second-order converses for the MAC, thus allowing us to obtain a conclusive second-order result for the Gaussian MAC with DMS.

We focus primarily on local second-order asymptotics propounded by Haim et al. [18] for general network information theory problems, where a boundary point is fixed and the rate of approach is characterized. This is different from the global asymptotics studied in [10]–[17], which we also study here as an initial step towards obtaining the local result. Similarly to Haim et al. [18], we believe that the study of local second-order asymptotics provides significantly greater insight into the system performance.

B. Main Contributions

Our main contribution is the characterization of the set of admissible second-order coding rates $(L_1, L_2)$ in the curved part of the boundary of the capacity region. For a point on the boundary characterized by $\rho \in (0, 1)$, we show that the set of achievable second-order rate pairs $(L_1, L_2)$ is given by those satisfying

$$\left[\frac{L_1}{L_1 + L_2}\right] \in \bigcup_{\beta \in \mathbb{R}} \left\{ \beta D(\rho) + \Psi^{-1}(V(\rho), \varepsilon) \right\},$$

where the entries of $D(\rho)$ are the derivatives of the capacities in (1)–(2), $V(\rho)$ is the dispersion matrix [10], [11], and $\Psi^{-1}$ is the 2-dimensional generalization of the inverse of the cumulative distribution of a Gaussian. (All quantities are defined precisely in the sequel.) Thus, the contribution from the Gaussian approximation $\Phi^{-1}(V(\rho), \varepsilon)$ is insufficient for characterizing the second-order asymptotics of multi-terminal channel coding problems in general, and we need to take into account contributions from the first-order term in terms of its slope $D(\rho)$. This is in stark
Fig. 2. Capacity region (CR) in nats/use of a Gaussian MAC with DMS where \( S_1 = S_2 = 1 \), i.e. 0 dB. Observe that \( \rho \in [0, 1] \) parametrizes points on the boundary. Every point on the curved part of the boundary is achieved by a unique input distribution \( \mathcal{N}(0, \Sigma(\rho)) \).

contrast to single-user problems (e.g. [3], [4], [6]–[8]) and the (two-encoder) Slepian-Wolf problem [5], [10] where the Gaussian approximation in terms of a dispersion quantity is sufficient for second-order asymptotics.

We first derive global second-order results [10], [18] and then use them to obtain local second-order results. As in [18], we make a strong connection between these two perspectives. To the best of our knowledge, our main result (Theorem 3) provides the first complete characterization of the second-order asymptotics of a multi-user information theory problem in which the boundary of the capacity region (or optimal rate region for source coding problems) is curved.

II. PROBLEM SETTING AND DEFINITIONS

In this section, we state the channel model, various definitions and some known results.

**Notation:** Given integers \( l \leq m \), we use the discrete interval notation \([l : m] := \{l, \ldots, m\} \) and \([m] := [1 : m] \). All \( \log \)’s and \( \exp \)’s are with respect to the natural base \( e \). The \( \ell_p \)-norm of the vectorized version of matrix \( \mathbf{A} \) is denoted by \( \| \mathbf{A} \|_p := \left( \sum_{i,j} |a_{i,j}|^p \right)^{1/p} \). For two vectors of the same length \( \mathbf{a}, \mathbf{b} \in \mathbb{R}^d \), the notation \( \mathbf{a} \leq \mathbf{b} \) means that \( a_j \leq b_j \) for all \( j \in [d] \). The notation \( \mathcal{N}(\mathbf{u}; \mu, \Lambda) \) denotes the multivariate Gaussian probability density function (pdf) with mean \( \mu \) and covariance \( \Lambda \). The argument \( \mathbf{u} \) will often remain unspecified. We use standard asymptotic notations: \( f_n \in O(g_n) \) if and only if (iff) \( \limsup_{n \to \infty} \left| f_n / g_n \right| < \infty \); \( f_n \in \Xi(g_n) \) iff \( g_n \in O(f_n) \); \( f_n \in \Theta(g_n) \) iff \( f_n \in O(g_n) \cap \Omega(g_n) \); \( f_n \in o(g_n) \) iff \( \limsup_{n \to \infty} \left| f_n / g_n \right| = 0 \); and \( f_n \in \omega(g_n) \) iff \( \liminf_{n \to \infty} \left| f_n / g_n \right| = \infty \).

A. Channel Model

The signal model for the Gaussian MAC is given by
\[
Y = X_1 + X_2 + Z, \tag{6}
\]
where $X_1$ and $X_2$ represent the inputs to the channel, $Z \sim N(0, 1)$ is additive Gaussian noise with zero mean and unit variance, and $Y$ is the output of the channel. Thus, the channel from $(X_1, X_2)$ to $Y$ can be written as

$$W(y|x_1, x_2) = \frac{1}{\sqrt{2\pi}} \exp\left(-\frac{1}{2}(y-x_1-x_2)^2\right).$$

(7)

The channel is used $n$ times in a memoryless manner without feedback. The channel inputs (i.e., the transmitted codewords) $x_1 = (x_{11}, \ldots, x_{1n})$ and $x_2 = (x_{21}, \ldots, x_{2n})$ are required to satisfy the maximal power constraints

$$\|x_1\|_2^2 \leq nS_1, \quad \text{and} \quad \|x_2\|_2^2 \leq nS_2,$$

(8)

where $S_1$ and $S_2$ are arbitrary positive numbers. We do not incorporate multiplicative gains $g_1$ and $g_2$ to $X_1$ and $X_2$ in the channel model in (6); this is without loss of generality, since in the presence of these gains we may equivalently redefine (8) with $S'_j := S_j/g_j^2$ for $j = 1, 2$.

**B. Definitions**

**Definition 1 (Code).** An $(n, M_{1,n}, M_{2,n}, S_1, S_2, \varepsilon_n)$-code for the Gaussian MAC with DMS consists of two encoders $f_{1,n}, f_{2,n}$, and a decoder $\varphi_n$ of the form $f_{1,n} : [M_{1,n}] \times [M_{2,n}] \rightarrow \mathbb{R}^n$, $f_{2,n} : [M_{2,n}] \rightarrow \mathbb{R}^n$ and $\varphi_n : \mathbb{R}^n \rightarrow [M_{1,n}] \times [M_{2,n}]$ which satisfy

$$\|f_{1,n}(m_1, m_2)\|_2^2 \leq nS_1 \quad \forall (m_1, m_2) \in [M_{1,n}] \times [M_{2,n}],$$

(9)

$$\|f_{2,n}(m_2)\|_2^2 \leq nS_2 \quad \forall m_2 \in [M_{2,n}],$$

(10)

$$\Pr((M_1, M_2) \neq (\hat{M}_1, \hat{M}_2)) \leq \varepsilon_n,$$

(11)

where the messages $M_1$ and $M_2$ are uniformly distributed on $[M_{1,n}]$ and $[M_{2,n}]$ respectively, and $(\hat{M}_1, \hat{M}_2) := \varphi_n(Y^n)$ are the decoded messages.

Since $S_1$ and $S_2$ are fixed positive numbers, we suppress the dependence of the subsequent definitions, results and parameters on these constants. We will often make reference to $(n, \varepsilon)$-codes; this is the family of $(n, M_{1,n}, M_{2,n}, S_1, S_2, \varepsilon)$-codes where the sizes $M_{1,n}, M_{2,n}$ are left unspecified.

**Definition 2 ((n, ε)-Achievability).** A pair of non-negative numbers $(R_1, R_2)$ is $(n, \varepsilon)$-achievable if there exists an $(n, M_{1,n}, M_{2,n}, S_1, S_2, \varepsilon_n)$-code such that

$$\frac{1}{n} \log M_{j,n} \geq R_j, \quad j = 1, 2, \quad \text{and} \quad \varepsilon_n \leq \varepsilon.$$

(12)

The $(n, \varepsilon)$-capacity region $\mathcal{C}(n, \varepsilon) \subset \mathbb{R}_+^2$ is defined to be the set of all $(n, \varepsilon)$-achievable rate pairs $(R_1, R_2)$.

Definition 2 is a non-asymptotic one that is used primarily for the global second-order results. We now introduce definitions that involve the existence of sequences of codes, implying that these are asymptotic-type definitions.

**Definition 3 (First-Order Coding Rates).** A pair of non-negative numbers $(R_1, R_2)$ is $\varepsilon$-achievable if there exists a sequence of $(n, M_{1,n}, M_{2,n}, S_1, S_2, \varepsilon_n)$-codes such that

$$\liminf_{n \to \infty} \frac{1}{n} \log M_{j,n} \geq R_j, \quad j = 1, 2, \quad \text{and} \quad \limsup_{n \to \infty} \varepsilon_n \leq \varepsilon.$$

(13)

The $\varepsilon$-capacity region $\mathcal{C}(\varepsilon) \subset \mathbb{R}_+^2$ is defined to be the closure of the set of all $\varepsilon$-achievable rate pairs $(R_1, R_2)$. The capacity-region $\mathcal{C}$ is defined as

$$\mathcal{C} := \bigcap_{\varepsilon > 0} \mathcal{C}(\varepsilon) = \lim_{\varepsilon \to 0} \mathcal{C}(\varepsilon),$$

(14)

where the limit exists because of the monotonicity of $\mathcal{C}(\varepsilon)$.

Next, we state the most important definitions concerning local second-order coding rates in the spirit of Nomura-Han [5] and Tan-Kosut [10]. We will spend the majority of the paper developing tools to characterize these rates. Here $(R_1^*, R_2^*)$ is a pair of rates on the boundary of $\mathcal{C}(\varepsilon)$. 
**Definition 4** (Second-Order Coding Rates). A pair of numbers \((L_1, L_2)\) is \((\varepsilon, R_1^*, R_2^*)\)-achievable if there exists a sequence of \((n, M_{1,n}, M_{2,n}, S_1, S_2, \varepsilon_n)\)-codes such that

\[
\liminf_{n \to \infty} \frac{1}{\sqrt{n}} (\log M_{j,n} - n R_j^*) \geq L_j, \quad j = 1, 2, \quad \text{and} \quad \limsup_{n \to \infty} \varepsilon_n \leq \varepsilon. \tag{15}
\]

The \((\varepsilon, R_1^*, R_2^*)\)-optimal second-order coding rate region \(\mathcal{L}(\varepsilon; R_1^*, R_2^*) \subset \mathbb{R}^2\) is defined to be the closure of the set of all \((\varepsilon, R_1^*, R_2^*)\)-achievable rate pairs \((L_1, L_2)\).

Stated differently, if \((L_1, L_2)\) is \((\varepsilon, R_1^*, R_2^*)\)-achievable, then there are codes whose error probabilities are asymptotically no larger than \(\varepsilon\), and whose sizes \((M_{1,n}, M_{2,n})\) satisfy the asymptotic relation in (3). Even though we refer to \(L_1\) and \(L_2\) as “rates”, they may be negative [3]–[6]. A negative value corresponds to a backoff from the first-order term, whereas a positive value corresponds to an addition to the first-order term.

**C. Existing First-Order Results**

To put things in context, let us review some existing results concerning the \(\varepsilon\)-capacity region. To state the result compactly, we define the mutual information vector as

\[
\mathbf{I}(\rho) = \begin{bmatrix} I_1(\rho) \\ I_{12}(\rho) \end{bmatrix} := \begin{bmatrix} C(S_1(1 - \rho^2)) \\ C(S_1 + S_2 + 2\rho\sqrt{S_1S_2}) \end{bmatrix}. \tag{16}
\]

For a pair of rates \((R_1, R_2)\), let the rate vector be

\[
\mathbf{R} := \begin{bmatrix} R_1 \\ R_1 + R_2 \end{bmatrix}. \tag{17}
\]

A statement of the following result is provided in [1] Ex. 5.18(b)]. A weak converse was proved for the more general Gaussian MAC with common message in [19].

**Proposition 1** (Capacity Region). The capacity region of the Gaussian MAC with DMS is given by

\[
\mathcal{C} = \bigcup_{0 \leq \rho \leq 1} \{(R_1, R_2) \in \mathbb{R}^2_+ : \mathbf{R} \leq \mathbf{I}(\rho)\}. \tag{18}
\]

The union on the right is a subset of \(\mathcal{C}(\varepsilon)\) for every \(\varepsilon \in (0, 1)\). However, only the weak converse is implied by (18). To the best of the authors’ knowledge, the strong converse for the Gaussian MAC with DMS has not been demonstrated. A by-product of the derivation of the second-order asymptotics in this paper is the strong converse, allowing us to assert that for all \(\varepsilon \in (0, 1)\),

\[
\mathcal{C} = \mathcal{C}(\varepsilon). \tag{19}
\]

The direct part of Proposition [1] can be proved using superposition coding [2]. Treat \(X_2\) as the cloud center and \(X_1\) as the satellite codeword. The input distribution to achieve a point on the boundary characterized by some \(\rho \in [0, 1]\) is a 2-dimensional Gaussian with zero mean and covariance matrix

\[
\Sigma(\rho) := \begin{bmatrix} S_1 & \rho\sqrt{S_1S_2} \\ \rho\sqrt{S_1S_2} & S_2 \end{bmatrix}. \tag{20}
\]

Thus, the parameter \(\rho\) represents the correlation between the two users’ codewords.

**III. GLOBAL SECOND-ORDER RESULTS**

In this section, we present inner and outer bounds to \(\mathcal{C}(n, \varepsilon)\). We begin with some definitions. Let \(V(x, y) := \frac{x(x+2)}{2(x+1)(y+1)}\) be the Gaussian cross-dispersion function and let \(V(x) := V(x, x)\) be the Gaussian dispersion function [4], [8], [9] for a single-user additive white Gaussian noise channel with signal-to-noise ratio \(x\). For fixed \(0 \leq \rho \leq 1\), define the information-dispersion matrix

\[
\mathbf{V}(\rho) := \begin{bmatrix} V_1(\rho) & V_{1,12}(\rho) \\ V_{1,12}(\rho) & V_{12}(\rho) \end{bmatrix}. \tag{21}
\]
Let \( \Phi \) be the Gaussian cdf of the standard normal distribution. These quantities can be thought of as the generalization of the cumulative distribution function (cdf) of the standard Gaussian \( \Phi(z) := \int_{-\infty}^{z} N(u; 0, 1) \, du \) and its inverse \( \Phi^{-1}(\varepsilon) := \sup \left\{ z \in \mathbb{R} : \Phi(z) \geq 1 - \varepsilon \right\} \) to the bivariate case.

\[
V(\rho) := \mathbf{V}(S_1(1 - \rho^2)) \geq 0, \\
V_{1,2}(\rho) := \mathbf{V}(S_1(1 - \rho^2), S_1 + S_2 + 2\rho\sqrt{S_1S_2}), \\
V_{2,2}(\rho) := \mathbf{V}(S_1 + S_2 + 2\rho\sqrt{S_1S_2}).
\]

Let \((X_1, X_2) \sim P_{X_1,X_2} = \mathcal{N}(0; \Sigma(\rho))\), and define \(Q_{Y|X_2}(y|x_2) := \mathcal{N}(y; x_2(1 + \rho\sqrt{S_1/S_2}), 1 + S_1(1 - \rho^2))\), \(Q_Y(y) := \mathcal{N}(y; 0, 1 + S_1 + S_2 + 2\rho\sqrt{S_1S_2})\).

\[
\begin{align*}
\mathbf{I}(\rho) &:= \mathbb{E}[\mathbf{j}(X_1, X_2, Y)], \\
\mathbf{V}(\rho) &:= \mathbb{E}[\text{Cov}(\mathbf{j}(X_1, X_2, Y) \mid X_1, X_2)].
\end{align*}
\]

For a given point \((z_1, z_2) \in \mathbb{R}^2\) and a (non-zero) positive semi-definite matrix \(\mathbf{V}\), define

\[
\Psi(z_1, z_2; \mathbf{V}) := \int_{-\infty}^{z_2} \int_{-\infty}^{z_1} N(u; 0, \mathbf{V}) \, du,
\]

and for a given \(\varepsilon \in (0, 1)\), define the set

\[
\Psi^{-1}(\mathbf{V}, \varepsilon) := \left\{ (z_1, z_2) \in \mathbb{R}^2 : \Psi(z_1, z_2; \mathbf{V}) \geq 1 - \varepsilon \right\}.
\]

These quantities can be thought of as the generalization of the cumulative distribution function (cdf) of the standard Gaussian \(\Phi(z) := \int_{-\infty}^{z} N(u; 0, 1) \, du\) and its inverse \(\Phi^{-1}(\varepsilon) := \sup \left\{ z \in \mathbb{R} : \Phi(z) \geq 1 - \varepsilon \right\}\) to the bivariate case.
We are now in a position to state our main result.

Let \( g(\rho, \varepsilon, n) \) and \( \overline{g}(\rho, \varepsilon, n) \) be arbitrary functions of \( \rho, \varepsilon \) and \( n \) for now, and define the inner and outer regions

\[
R_{\text{in}}(n, \varepsilon; \rho) := I(\rho) + \frac{\Psi^{-1}(V(\rho), \varepsilon)}{\sqrt{n}} + g(\rho, \varepsilon, n) 1,
\]

\[
R_{\text{out}}(n, \varepsilon; \rho) := I(\rho) + \frac{\Psi^{-1}(V(\rho), \varepsilon)}{\sqrt{n}} + \overline{g}(\rho, \varepsilon, n) 1.
\]

**Theorem 2** (Global Bounds on the \((n, \varepsilon)\)-Capacity Region). There exist functions \( g(\rho, \varepsilon, n) \) and \( \overline{g}(\rho, \varepsilon, n) \) such that the \((n, \varepsilon)\)-capacity region satisfies

\[
\bigcup_{0 \leq \rho \leq 1} R_{\text{in}}(n, \varepsilon; \rho) \subset C(n, \varepsilon) \subset \bigcup_{-1 \leq \rho \leq 1} R_{\text{out}}(n, \varepsilon; \rho),
\]

and such that \( g \) and \( \overline{g} \) satisfy the following properties:

1) For any \( \varepsilon \in (0, 1) \) and \( \rho \in (-1, 1) \), we have

\[
g(\rho, \varepsilon, n) = O\left(\frac{\log n}{n}\right), \quad \text{and} \quad \overline{g}(\rho, \varepsilon, n) = O\left(\frac{\log n}{n}\right).
\]

2) For any \( \varepsilon \in (0, 1) \) and any sequence \( \{\rho_n\} \) with \( \rho_n \to \pm 1 \), we have

\[
g(\rho_n, \varepsilon, n) = o\left(\frac{1}{\sqrt{n}}\right), \quad \text{and} \quad \overline{g}(\rho_n, \varepsilon, n) = o\left(\frac{1}{\sqrt{n}}\right).
\]

The proof of Theorem 2 is provided in Section VI. We remark that even though the union for the outer bound is taken over \( \rho \in [-1, 1] \), only the values \( \rho \in [0, 1] \) will play a role in establishing the local asymptotics in Section IV since negative values of \( \rho \) are not even first-order optimal, i.e. they fail to achieve a point on the boundary of the capacity region.

**IV. LOCAL SECOND-ORDER CODING RATES**

In this section, we present our main result, namely, the characterization of the \((\varepsilon, R_1^*, R_2^*)\)-optimal second-order coding rate region \( \mathcal{L}(\varepsilon; R_1^*, R_2^*) \) (see Definition 4), where \((R_1^*, R_2^*) \) is an arbitrary point on the boundary of \( \mathcal{C} \). Our result is stated in terms of the derivative of the mutual information vector with respect to \( \rho \), namely

\[
D(\rho) = \begin{bmatrix} D_1(\rho) \\ D_2(\rho) \end{bmatrix} := \frac{\partial}{\partial \rho} \begin{bmatrix} I_1(\rho) \\ I_2(\rho) \end{bmatrix},
\]

where the individual derivatives are given by

\[
\frac{\partial I_1(\rho)}{\partial \rho} = \frac{-S_1 \rho}{1 + S_1 (1 - \rho^2)} \quad \text{(38)}
\]

\[
\frac{\partial I_2(\rho)}{\partial \rho} = \frac{\sqrt{S_1 S_2}}{1 + S_1 + S_2 + 2 \rho \sqrt{S_1 S_2}} \quad \text{(39)}
\]

Note that \( \rho \in (0, 1) \) represents the strictly concave part of the boundary (see Fig. 2), and in this interval we have \( D_1(\rho) < 0 \) and \( D_1(\rho) > 0 \).

We introduce the following notation: For a vector \( \mathbf{v} = (v_1, v_2) \in \mathbb{R}^2 \), define the down-set of \( \mathbf{v} \) as

\[
\mathbf{v}^- := \{(w_1, w_2) \in \mathbb{R}^2 : w_1 \leq v_1, w_2 \leq v_2\}.
\]

We are now in a position to state our main result.

**Theorem 3** (Optimal Second-Order Coding Rate Region). Depending on \((R_1^*, R_2^*)\), we have the following three cases:

(i) If \( R_1^* = I_1(0) \) and \( R_1^* + R_2^* \leq I_2(0) \) (vertical segment of the boundary corresponding to \( \rho = 0 \)), then

\[
\mathcal{L}(\varepsilon; R_1^*, R_2^*) = \left\{(L_1, L_2) \in \mathbb{R}^2 : L_1 \leq \sqrt{V_1(0) \Phi^{-1}(\varepsilon)} \right\}.
\]

\[
\Phi^{-1}(\varepsilon)
\]

\[
\sqrt{V_1(0) \Phi^{-1}(\varepsilon)}
\]

\[
\left\{(L_1, L_2) \in \mathbb{R}^2 : L_1 \leq \sqrt{V_1(0) \Phi^{-1}(\varepsilon)} \right\}.
\]
(ii) If \( R_1^* = I_1(\rho) \) and \( R_1^* + R_2^* = I_{12}(\rho) \) (curved segment of the boundary corresponding to \( 0 < \rho < 1 \)), then
\[
\mathcal{L}(\varepsilon; R_1^*, R_2^*) = \left\{ (L_1, L_2) \in \mathbb{R}^2 : \left[ \frac{L_1}{L_1 + L_2} \right] \in \bigcup_{\beta \in \mathbb{R}} \left\{ \beta D(\rho) + \Psi^{-1}(V(\rho), \varepsilon) \right\} \right\}. \tag{42}
\]

(iii) If \( R_1^* = 0 \) and \( R_1^* + R_2^* = I_{12}(1) \) (point on the vertical axis corresponding to \( \rho = 1 \)), then
\[
\mathcal{L}(\varepsilon; R_1^*, R_2^*) = \left\{ (L_1, L_2) \in \mathbb{R}^2 : \left[ \frac{L_1}{L_1 + L_2} \right] \in \bigcup_{\beta \leq 0} \left\{ \beta D(1) + \left[ \frac{0}{\sqrt{V_{12}(1)\Psi^{-1}(\varepsilon)}} \right] \right\} \right\}. \tag{43}
\]

The proof of Theorem 3 is provided in Section VII. It leverages on the global second-order result in Theorem 2.

A. Discussion

Observe that in case (i), the second-order region is simply characterized by a scalar dispersion term \( V_1(0) \) and the inverse of the Gaussian cdf \( \Phi^{-1} \). Roughly speaking, in this part of the boundary, there is effectively only a single rate constraint in terms of \( R_1 \), since we are operating “far away” from the sum rate constraint. This results in a large deviations-type event for the sum rate constraint which has no bearing on second-order asymptotics; see further discussions in [5], [10] and [18].

Cases (ii)–(iii) are more interesting, and their proofs are non-trivial. As in Nomura-Han [5] and Tan-Kosut [10], the second-order asymptotics for case (ii) depend on the dispersion matrix \( V(\rho) \) and the 2-dimensional analogue of the inverse of the Gaussian cdf \( \Psi^{-1} \), since both rate constraints are active at a point on the boundary parametrized by \( \rho \in (0, 1) \). However, in our setting, the expression containing \( \Psi^{-1} \) alone (i.e. the expression obtained by setting \( \beta = 0 \) in (42)) corresponds to only considering the unique input distribution \( N(0, \Sigma(\rho)) \) achieving the point \((I_1(\rho), I_{12}(\rho) - I_1(\rho))\). From Fig. 2, this is not sufficient to achieve all second-order coding rates, since there are non-empty regions within the capacity region that are not contained in the trapezium of rate pairs achievable using \( N(0, \Sigma(\rho)) \).

Thus, to achieve all \((L_1, L_2)\) pairs, we must allow the input distribution to vary with the blocklength \( n \). This is manifested in the \( \beta D(\rho) \) term. Our proof of the direct part involves random coding with an input distribution of the form \( N(0, \Sigma(\rho_n)) \), where \( \rho_n - \rho = O\left( \frac{1}{\sqrt{n}} \right) \). By a Taylor series, the resulting mutual information vector \( I(\rho_n) \) is approximately \( I(\rho) + (\rho_n - \rho)D(\rho) \). Since \( \rho_n - \rho = O\left( \frac{1}{\sqrt{n}} \right) \), the gradient term \( (\rho_n - \rho)D(\rho) \) also contributes to the second-order behavior, together with the Gaussian approximation term \( \Psi^{-1}(V(\rho), \varepsilon) \).

For the converse, we consider an arbitrary sequence of codes with rate pairs \( \{(R_{1,n}, R_{2,n})\}_{n \geq 1} \) converging to \((I_1(\rho), I_{12}(\rho) - I_1(\rho))\) with second-order behavior given by (15). From the global result, we know \( [R_{1,n}, R_{1,n} + R_{2,n}] \in \mathcal{R}_{out}(n, \varepsilon; \rho_n) \) for some sequence of \( \rho_n \). We then establish, using the definition of the second-order coding rates in (15), that \( \rho_n = \rho + O\left( \frac{1}{\sqrt{n}} \right) \). Finally, by the Bolzano-Weierstrass theorem [20, Thm. 3.6(b)], we may pass to a subsequence of \( \rho_n \) (if necessary), and this establishes the converse.

A similar discussion holds true for case (iii); the main differences are that the covariance matrix is singular, and that the union in (43) is taken over \( \beta \leq 0 \) only, since \( \rho_n \) can only approach one from below.

B. Second-Order Asymptotics for a Given Angle of Approach

Here we study the second-order behavior when a point on the boundary is approached from a given angle, as was done in Tan-Kosut [10]. If \( \varepsilon < \frac{1}{2} \) (resp. \( \varepsilon > \frac{1}{2} \)), we approach a boundary point from inside (resp. outside) the capacity region. We focus on the most interesting case in Theorem 3, namely, case (ii) corresponding to \( \rho \in (0, 1) \). Case (iii) can be handled similarly, and in case (i) the angle of approach is of little interest since \( L_2 \) can be arbitrarily large or small.

First, we present an alternative expression for the set \( \mathcal{L} = \mathcal{L}(\varepsilon; R_1^*, R_2^*) \) given in (42) with \( R_1^* = I_1(\rho) \) and \( R_1^* + R_2^* = I_{12}(\rho) \) for some \( \rho \in (0, 1) \). It is easily seen that \( (L_1, L_2) \in \mathcal{L} \) implies \( (L_1 + \beta D_1(\rho), L_2 + \beta D_2(\rho)) \in \mathcal{L} \), where \( D_2(\rho) := D_{12}(\rho) - D_1(\rho) \). It follows that \( \mathcal{L} \) equals the set of all points lying below a straight line with slope \( \frac{D_2(\rho)}{D_1(\rho)} \), which intersects the boundary of \( \Psi^{-1}(V(\rho), \varepsilon) \). That is,
\[
\mathcal{L}(\varepsilon; R_1^*, R_2^*) = \left\{ (L_1, L_2) : L_2 \leq a_\rho L_1 + b_{\rho, \varepsilon} \right\}, \tag{44}
\]
For which coding with a fixed (not varying with $\varepsilon$) $\beta$, there is only one angle $\theta$ for which the intersection of the boundary of $\mathcal{L}(\varepsilon; R_1^*, R_2^*)$ and the line $L_2 = L_1 \tan \theta$ coincides with the boundary of $\Psi^{-1}(\mathbf{V}(\rho), \varepsilon)$.

where

$$a_\rho := \frac{D_2(\rho)}{D_1(\rho)}, \quad \text{and} \quad b_{\rho, \varepsilon} := \inf \left\{ b : \left( L_1, a_\rho L_1 + b \right) \in \Psi^{-1}(\mathbf{V}(\rho), \varepsilon) \text{ for some } L_1 \in \mathbb{R} \right\}. \quad (45)$$

We provide an example in Fig. 4 with the parameters $S_1 = S_2 = 1$, $\rho = \frac{1}{2}$ and $\varepsilon = 0.1$. Since $\varepsilon < \frac{1}{2}$, the boundary point $(R_1^*, R_2^*)$ is approached from the inside. See Fig. 3 where for $\varepsilon < \frac{1}{2}$, the set $\Psi^{-1}(\mathbf{V}, \varepsilon)$ only contains points with negative coordinates.

Given the gradient $a_\rho$, the offset $b_{\rho, \varepsilon}$, and an angle $\theta$ (measured with respect to the horizontal axis), we seek the pair $(L_1, L_2)$ on the boundary of $\mathcal{L}(\varepsilon; R_1^*, R_2^*)$ such that $L_2 = L_1 \tan \theta$. It is easily seen that this point is obtained by solving for the intersection of the line $L_2 = a_\rho L_1 + b_{\rho, \varepsilon}$ with $L_2 = L_1 \tan \theta$. The two lines coincide when

$$L_1 = \frac{b_{\rho, \varepsilon}}{\tan \theta - a_\rho}, \quad \text{and} \quad L_2 = \frac{b_{\rho, \varepsilon} \tan \theta}{\tan \theta - a_\rho}. \quad (46)$$

In Fig. 4 we see that there is only a single angle $\theta_{\rho, \varepsilon}^* \approx 3.253$ rads for which the point of intersection in (46) is also on the boundary of $\Psi^{-1}(\mathbf{V}(\rho), \varepsilon)$, yielding $(L_1, L_2) \approx (-0.920, -0.103)$. In other words, there is only one angle for which coding with a fixed (not varying with $n$) input distribution $\mathcal{N}(0, \mathbf{V}(\rho))$ is optimal in the second-order sense (i.e. for which the added term $\beta \mathbf{D}(\rho)$ in (42) is of no additional help and $\beta = 0$ is optimal). For all the other angles, we should choose a non-zero coefficient $\beta$, which corresponds to choosing an input distribution which varies with $n$.

Finally, in Fig. 5 we plot the norm of the vector of second-order rates $[L_1, L_2]^T$ in (46) against $\theta$, the angle of approach. For $\varepsilon < \frac{1}{2}$, the point $[L_1, L_2]^T$ may be interpreted as that corresponding to the “smallest backoff” from...
Fig. 5. Plot of $\sqrt{L_1^2 + L_2^2}$ against the angle of approach $\theta$ with the same parameters as in Fig. 4. The second-order rates $L_1, L_2$, as functions of $\theta$, are defined in [46]. Here, $\pi + \arctan(a_\rho) \approx 0.328 \times (2\pi)$ and $2\pi + \arctan(a_\rho) \approx 0.828 \times (2\pi)$ are the critical angles where $\sqrt{L_1^2 + L_2^2}$ diverges.

Thus, $\sqrt{L_1^2 + L_2^2}$ is a measure of the total backoff. For $\varepsilon > \frac{1}{2}$, $[L_1, L_2]^T$ corresponds to the “largest addition” to the first-order rates. It is noted that the norm tends to infinity when the angle tends to $\pi + \arctan(a_\rho)$ (from above) or $2\pi + \arctan(a_\rho)$ (from below). This corresponds to an approach almost parallel to the gradient at the point on the boundary parametrized by $\rho$. A similar phenomenon was observed for the Slepian-Wolf problem [10].

V. CONCLUDING REMARKS

We have studied the second-order asymptotics (i.e. identified the optimal second-order coding rate region) of the Gaussian MAC with DMS. There are two reasons as to why the analysis here is more tractable vis-à-vis finite blocklength or second-order analysis for the the discrete memoryless MAC (DM-MAC) studied extensively in [10]–[12], [16]–[18]. Gaussianity allows us to identify the boundary of the capacity region and associate each point on the boundary with an input distribution parametrized by $\rho$. Because for the DM-MAC, one needs to take the convex closure of the union over input distributions $P_{X_1, X_2}$ to define the capacity region [1, Sec. 4.5], the boundary points are more difficult to characterize. In addition, in the absence of the DMS assumption, one needs to ensure in a converse proof (possibly related to the wringing technique of Ahlswede [21]) that the codewords pairs are almost orthogonal. By leveraging on the DMS assumption, we circumvent this requirement.

For future investigations, we note that the Gaussian broadcast channel [1, Sec. 5.5] is a problem which is very similar to the Gaussian MAC with DMS (both require superposition coding and each point on the boundary is achieved by a unique input distribution). As such, we expect that some of the second-order analysis techniques contained herein may be applicable to the Gaussian broadcast channel.

1There may be some imprecision in the use of the word “backoff” here as for angles in the second (resp. fourth) quadrant, $L_2$ (resp. $L_1$) is positive. On the other hand, one could generally refer to “backoff” as moving in some inward direction (relative to the capacity region boundary) even if it is in a direction where one of the second-order rates increases. The same goes for the term “addition”.
VI. PROOF OF THEOREM 2: GLOBAL SECOND-ORDER RESULT

A. Converse Part

We first prove the outer bound in (34). The analysis is split into six steps.

1) A Reduction from Maximal to Equal Power Constraints: We first perform a reduction that will simplify the proof. Let \( C_{eq}(n, \varepsilon) \) be the \((n, \varepsilon)\)-capacity region in the case that \((9)\) and \((10)\) are equality constraints, i.e., \( \|f_{1,n}(m_1, m_2)\|_2 = nS_1 \) and \( \|f_{2,n}(m_2)\|_2 = nS_2 \) for all \((m_1, m_2)\). We claim that
\[
C_{eq}(n, \varepsilon) \subset C(n, \varepsilon) \subset C_{eq}(n + 1, \varepsilon). \tag{47}
\]
The lower bound is obvious. The upper bound follows by noting that the decoder for the length-\((n + 1)\) code can ignore the last symbol, which can be chosen to equalize the powers.

It follows from \((47)\) that for the purpose of second-order asymptotics, \( C_{eq}(n, \varepsilon) \) and \( C(n, \varepsilon) \) are equivalent. This argument was also used in \([8, \text{Lem. 39}]\) and \([9, \text{Sec. XIII}]\). Henceforth, we assume that all codewords \((x_1, x_2)\) have empirical powers exactly equal to \((S_1, S_2)\).

2) Correlation Type Classes: Define \( I_0 := \{0\} \) and \( I_k := \left(\frac{k-1}{n}, \frac{k}{n}\right], k \in [n] \), and let \( I_{-k} := -I_k \) for \( k \in [n] \). We see that the family \( \{I_k : k \in [-n : n]\} \) forms a partition of \([-1, 1]\). Consider the correlation type classes (or simply type classes)
\[
T_n(k) := \left\{ (x_1, x_2) : \frac{\langle x_1, x_2 \rangle}{\|x_1\|_2 \|x_2\|_2} \in I_k \right\} \tag{48}
\]
where \( k \in [-n : n] \), and \( \langle x_1, x_2 \rangle := \sum_{i=1}^{n} x_{1i} x_{2i} \) is the standard inner product in \( \mathbb{R}^n \). The total number of type classes is \( 2n + 1 \), which is polynomial in \( n \) analogously to the case of discrete alphabets \([22, \text{Ch. 2}]\).

We use an argument from Csiszár-Körner \([22, \text{Lem. 16.2}]\) to perform a further reduction (along with the one in Section VI (A1)) to codes for which all codeword pairs have the same type. Let the message sets of the general code be \( M_{j,n} = [M_{j,n}] \) for \( j = 1, 2 \). Let a codebook \( C := \{(x_1(m_1, m_2), x_2(m_2)) : m_1 \in M_{1,n}, m_2 \in M_{2,n}\} \) be given. For each \( m_2 \in M_{2,n} \), we can find a set \( \tilde{M}_{1,n}(m_2) \subset M_{1,n} \) such that all pairs of codewords \((x_1(m_1, m_2), x_2(m_2))\), \( m_1 \in \tilde{M}_{1,n}(m_2) \) have the same type, say indexed by \( k(m_2) \in [-n : n] \), and
\[
\frac{1}{n} \log |\tilde{M}_{1,n}(m_2)| \geq \frac{1}{n} \log |M_{1,n}(m_2)| - \frac{\log(2n + 1)}{n}, \quad \forall m_2 \in M_{2,n}. \tag{49}
\]
We may assume that all the sets \( \tilde{M}_{1,n}(m_2), m_2 \in M_{2,n} \) have the same cardinality; otherwise, we can remove extra codeword pairs from some sets \( \tilde{M}_{1,n}(m_2) \) and \((49)\) will still be satisfied. We may also assume (by relabeling if necessary) that \( M_{1,n} := M_{1,n}(m_2) \) is the same set for all \( m_2 \). Now, we have a subcodebook \( \tilde{C}_1 := \{(x_1(m_1, m_2), x_2(m_2)) : m_1 \in \tilde{M}_{1,n}, m_2 \in M_{2,n}\} \), where for each \( m_2 \), all the codeword pairs have the same type and \((49)\) is satisfied. Across the \( m_2 \)'s, there may be different types indexed by \( k(m_2) \in [-n : n] \), but there exists a dominant type indexed by \( k^* \in \{k(m_2) : m_2 \in M_{2,n}\} \) and a set \( \tilde{M}_{2,n} \subset M_{2,n} \) such that
\[
\frac{1}{n} \log |\tilde{M}_{2,n}| \geq \frac{1}{n} \log |M_{2,n}| - \frac{\log(2n + 1)}{n}. \tag{50}
\]
As such, we have shown that there exists a subcodebook \( \tilde{C}_{12} := \{(x_1(m_1, m_2), x_2(m_2)) : m_1 \in \tilde{M}_{1,n}, m_2 \in \tilde{M}_{2,n}\} \) of constant type indexed by \( k^* \) where the sum rate satisfies
\[
\frac{1}{n} \log |\tilde{M}_{1,n} \times \tilde{M}_{2,n}| \geq \frac{1}{n} \log |M_{1,n} \times M_{2,n}| - \frac{2\log(2n + 1)}{n}. \tag{51}
\]
Combining \((50)\) and \((51)\), we see that the converse part of Theorem 2 for fixed-type codes implies the same for general codes, since the additional \( O\left(\frac{\log(2n)}{n}\right) \) factors in \((50)\) and \((51)\) can be absorbed into the third-order term \( \tilde{g}(\rho, \varepsilon, n) \). Thus, in the remainder of the proof, we limit our attention to fixed-type codes. For each \( n \), the type is indexed by \( k \in [-n : n] \), and we define \( \hat{\rho} := \frac{k}{n} \in [-1, 1] \). In some cases, we will be interested in \textit{sequences} of such values, in which case we will make the dependence on \( n \) explicit by writing \( \hat{\rho}_n \).
3) A Verdù-Han-type Converse Bound for the Gaussian MAC with DMS: We now state a non-asymptotic converse bound for the Gaussian MAC with DMS based on analogous bounds in Han’s work on the information spectrum approach for the general MAC [23, Lem. 4] and in Boucheron-Salamatian’s work on the information spectrum approach for the general broadcast channel with DMS [24, Lem. 2]. This result can be proved similarly to [23, 24], so we omit its proof.

**Proposition 4.** Fix a blocklength \( n \geq 1 \), auxiliary output distributions \( Q_{Y|X_2} \) and \( Q_Y \), and a constant \( \gamma > 0 \). For any \((n, M_{1,n}, M_{2,n}, S_1, S_2, \varepsilon_n)\)-code with codewords of fixed empirical powers \( S_1 \) and \( S_2 \) falling into a single correlation type class \( T_n(k) \), there exist random vectors \((X_1, X_2)\) with joint distribution \( P_{X_1,X_2} \) supported on \( \{(x_1, x_2) \in T_n(k) : \|x_j\|_2^2 = nS_j, j = 1, 2\} \) such that

\[
\varepsilon_n \geq \Pr(A_n \cup B_n) - 2\exp(-n\gamma), \tag{52}
\]

where

\[
A_n := \left\{ \frac{1}{n} \log \frac{W^n(Y|X_1, X_2)}{Q_{Y|X_2}(Y|X_2)} \leq \frac{1}{n} \log M_{1,n} - \gamma \right\},
\]

\[
B_n := \left\{ \frac{1}{n} \log \frac{W^n(Y|X_1, X_2)}{Q_Y(Y)} \leq \frac{1}{n} \log \left( M_{1,n} M_{2,n} \right) - \gamma \right\},
\]

with \( Y \mid \{X_1 = x_1, X_2 = x_2\} \sim W^n(\cdot|x_1, x_2) \).

There are several differences in Proposition 4 compared to [23, Lem. 4]. First, in our work, there are cost constraints on the codewords, and thus the support of the input distribution \( P_{X_1,X_2} \) is specified to reflect this constraint. Second, there are two (instead of three) events in the probability in (52) because the informed encoder \( f_{1,n} \) has access to both messages. Third, we can choose arbitrary output distributions \( Q_{Y|X_2} \) and \( Q_Y \). This generalization is analogous to the non-asymptotic converse bound by Hayashi and Nagaoka for classical-quantum channels [25, Lem. 4]. The freedom to choose the output distribution is crucial in both our problem and [25].

4) Evaluation of the Verdù-Han Bound for \( \hat{\rho} \in (-1, 1) \): Recall from Sections VI-A1 and VI-A2 that the codewords satisfy exact power constraints and belong to a single type class \( T_n(k) \). In this subsection, we consider the case that \( \hat{\rho} := \frac{k}{n} \in (-1, 1) \), and we derive bounds which will be useful for sequences \( \hat{\rho}_n \) uniformly bounded away from \(-1\) and \( 1 \). In Section VI-A3, we present alternative bounds to handle the remaining cases.

We evaluate (52) for a single correlation type class parametrized by \( k \). Let \( \gamma := \frac{\log n}{2n} \), yielding \( 2\exp(-n\gamma) = \frac{2}{\sqrt{n}} \).

We choose the output distributions \( Q_{Y|X_2} \) and \( Q_Y \) to be the \( n \)-fold products of \( Q_{Y|X_2} \) and \( Q_Y \), defined in (25)–(26) respectively, with \( \hat{\rho} \) in place of \( \rho \).

We now characterize the statistics of the first and second moments of \( j(x_{1i}, x_{2i}, Y_i) \) in (27) for fixed sequences \((x_1, x_2) \in T_n(k) \). From Appendix A these moments can be expressed as continuously differentiable functions of the empirical powers \( \frac{1}{n}\|x_1\|_2^2, \frac{1}{n}\|x_2\|_2^2 \) and the empirical correlation coefficient \( \frac{\langle x_1, x_2 \rangle}{\|x_1\|_2 \|x_2\|_2} \). The former two quantities are fixed due to the reduction in Section VI-A1 and the latter is within \( \frac{1}{n} \) of \( \hat{\rho} \) by the assumption that \((x_1, x_2) \in T_n(k) \). Thus, Taylor expanding (A.6) and (A.12) in Appendix A we obtain

\[
\left\| \mathbb{E} \left[ \frac{1}{n} \sum_{i=1}^{n} j(x_{1i}, x_{2i}, Y_i) \right] - I(\hat{\rho}) \right\|_{\infty} \leq \frac{\xi_1}{n}, \tag{55}
\]

\[
\left\| \text{Cov} \left[ \frac{1}{\sqrt{n}} \sum_{i=1}^{n} j(x_{1i}, x_{2i}, Y_i) \right] - V(\hat{\rho}) \right\|_{\infty} \leq \frac{\xi_2}{n}, \tag{56}
\]

for some \( \xi_1 > 0 \) and \( \xi_2 > 0 \) which can taken to be independent of \( \hat{\rho} \) (since the corresponding derivatives are uniformly bounded).

Let \( R_{j,n} := \frac{1}{n} \log M_{j,n} \) for \( j = 1, 2 \), and let \( R_n := [R_{1,n}, R_{1,n} + R_{2,n}]^T \). We have

\[
\Pr(A_n \cup B_n) = 1 - \Pr(A_n^c \cap B_n^c) = 1 - \mathbb{E}_{X_1, X_2} \left[ \Pr(A_n^c \cap B_n^c | X_1, X_2) \right] \tag{57}
\]
and in particular, using the definition of \( j(x_1, x_2, y) \) in (27) and the fact that \( Q_{Y|X_2} \) and \( Q_Y \) are product distributions,

\[
\Pr(A^c_n \cap B^c_n | x_1, x_2) = \Pr \left( \frac{1}{n} \sum_{i=1}^{n} j(x_{1i}, x_{2i}, Y_i) > R_n - \gamma 1 \right)
\]

(58)

\[
\leq \Pr \left( \frac{1}{n} \sum_{i=1}^{n} \left( j(x_{1i}, x_{2i}, Y_i) - \mathbb{E}[j(x_{1i}, x_{2i}, Y_i)] \right) > R_n - I(\hat{\rho}) - \gamma 1 - \frac{\xi_1}{n} 1 \right),
\]

(59)

where (59) follows from (55).

We are now in a position to apply the multivariate Berry-Esseen theorem [26], [27] (see Appendix B). The first two moments are bounded according to (55)–(56), and in Appendix A we show that, upon replacing the given \((x_1, x_2)\) pair with a different pair yielding the same statistics of \( \sum_{i=1}^{n} j(x_{1i}, x_{2i}, Y_i) \) if necessary (cf. Lemma 9), the required third moment is uniformly bounded (cf. Lemma 10). It follows that

\[
\Pr(A^c_n \cap B^c_n | x_1, x_2) 
\leq \Psi \left( \sqrt{n} \left( \frac{1}{n} \sum_{i=1}^{n} j(x_{1i}, x_{2i}, Y_i) \right) \right) + \frac{\psi(\hat{\rho}) \log n}{\sqrt{n}}.
\]

(60)

By Taylor expanding the continuously differentiable function \( \Psi(z_1, z_2, V) \mapsto \Psi(z_1, z_2; V) \), using the approximation in (56) and the fact that \( \det(V(\hat{\rho})) > 0 \) for \( \hat{\rho} \in (-1, 1) \), we obtain

\[
\Pr(A^c_n \cap B^c_n | x_1, x_2) \leq \Psi \left( \sqrt{n} \left( I_1(\hat{\rho}) - R_{1n} \right), \sqrt{n} \left( I_2(\hat{\rho}) - (R_{1n} + R_{2n}) \right) ; V(\hat{\rho}) \right) + \frac{\eta(\hat{\rho}) \log n}{\sqrt{n}}.
\]

(61)

The following “continuity” lemma for \( \varepsilon \mapsto \Psi^{-1}(V, \varepsilon) \) is proved in Appendix C

**Lemma 5.** Fix \( 0 < \varepsilon < 1 \) and a positive sequence \( \lambda_n = o(1) \). Let \( V \) be a non-zero positive semi-definite matrix. There exists a function \( h(V, \varepsilon) \) such that

\[
\Psi^{-1}(V, \varepsilon + \lambda_n) \subset \Psi^{-1}(V, \varepsilon) + h(V, \varepsilon) \lambda_n 1.
\]

(63)

We conclude from Lemma 5 that

\[
\Psi^{-1}\left( V(\hat{\rho}), \varepsilon + \frac{2}{\sqrt{n}} + \frac{\eta(\hat{\rho}) \log n}{\sqrt{n}} \right) \subset \Psi^{-1}(V(\hat{\rho}), \varepsilon) + h(\hat{\rho}, \varepsilon) \log n 1
\]

(64)

for some function \( h(\hat{\rho}, \varepsilon) := h(V(\hat{\rho}), \varepsilon) \) which diverges only as \( \hat{\rho} \to \pm 1 \). Uniting (62) and (64), we deduce that

\[
\left[ \frac{R_{1n}}{R_{1n} + R_{2n}} \right] \in I(\hat{\rho}) + \frac{\Psi^{-1}(V(\hat{\rho}), \varepsilon) + h(\hat{\rho}, \varepsilon) \log n}{\sqrt{n}} 1.
\]

(65)

5) **Evaluation of the Verdú-Han Bound with \( \hat{\rho}_n \to \pm 1 \):** Here we consider a sequence of codes of a single type indexed by \( k_n \) such that \( \hat{\rho}_n := \frac{k_n}{n} \to 1 \). The case \( \hat{\rho}_n \to -1 \) is handled similarly, and the details are thus omitted. Our aim is to show that

\[
\left[ \frac{R_{1n}}{R_{1n} + R_{2n}} \right] \in I(\hat{\rho}_n) + \frac{\Psi^{-1}(V(\hat{\rho}_n), \varepsilon)}{\sqrt{n}} + o\left( \frac{1}{\sqrt{n}} \right) 1.
\]

(66)

The following lemma states that as \( \hat{\rho}_n \to 1 \), the set \( \Psi^{-1}(V(1), \varepsilon) \) in (66) can be approximated by \( \Psi^{-1}(V(1), \varepsilon) \), which is a simpler rectangular set. The proof of the lemma is provided in Appendix D.
Lemma 6. Fix \( 0 < \varepsilon < 1 \) and let \( \hat{\rho}_n \to 1 \). There exist positive sequences \( a_n, b_n = \Theta((1 - \hat{\rho}_n)^{1/4}) \) and \( c_n = \Theta((1 - \hat{\rho}_n)^{1/2}) \) satisfying

\[
\left[ \sqrt{\frac{1}{V_{12}(1)}} \Phi^{-1}(\varepsilon + a_n) \right] - b_n \mathbf{1} \subset \Psi^{-1}(V(\hat{\rho}_n), \varepsilon) \subset \left[ \sqrt{\frac{1}{V_{12}(1)}} \Phi^{-1}(\varepsilon) \right] - c_n \mathbf{1}.
\] (67)

The down-set notation \([v]^-\) is defined in (40).

From the inner bound in Lemma 6, in order to show (66) it suffices to show

\[
\begin{bmatrix} R_{1,n} \\ R_{1,n} + R_{2,n} \end{bmatrix} \leq \mathbf{I}(\hat{\rho}_n) + \sqrt{\frac{1}{V_{12}(1)}} \begin{bmatrix} 0 \\ \Phi^{-1}(\varepsilon) \end{bmatrix} + o\left(\frac{1}{n}\right) \mathbf{1},
\] (68)

where we absorbed the sequences \( a_n, b_n \) into the \( o\left(\frac{1}{n}\right) \) term.

We return to the step in (59), which when combined with the Verdú-Han bound (with \( \gamma := \frac{\log n}{2n} \)) yields for some \((x_1, x_2) \in T_n(k)\) that

\[
\varepsilon_n \geq 1 - \Pr \left( \frac{1}{n} \sum_{i=1}^{n} \left( j(x_{1i}, x_{2i}, Y_i) - \mathbb{E}[j(x_{1i}, x_{2i}, Y_i)] \right) > R_n - \mathbf{I}(\hat{\rho}_n) - \gamma - \frac{\xi_1}{n} \right) - \frac{2}{\sqrt{n}}
\] (69)

\[
\geq \max \left\{ \Pr \left( \frac{1}{n} \sum_{i=1}^{n} \left( j_1(x_{1i}, x_{2i}, Y_i) - \mathbb{E}[j_1(x_{1i}, x_{2i}, Y_i)] \right) \leq R_{1,n} - I_1(\hat{\rho}_n) - \gamma - \frac{\xi_1}{n} \right) ,
\right.
\]

\[
\Pr \left( \frac{1}{n} \sum_{i=1}^{n} \left( j_2(x_{1i}, x_{2i}, Y_i) - \mathbb{E}[j_2(x_{1i}, x_{2i}, Y_i)] \right) \leq R_{1,n} + R_{2,n} - I_1(\hat{\rho}_n) - \gamma - \frac{\xi_1}{n} \right) \right\} - \frac{2}{\sqrt{n}}.
\] (70)

From (56) and the assumption that \( \hat{\rho}_n \to 1 \), the variance of \( \sum_{i=1}^{n} j_2(x_{1i}, x_{2i}, Y_i) \) equals \( n(V_{12}(1) + o(1)) \). Since \( V_{12}(1) > 0 \), we can treat the second term in the maximum in (70) in an identical fashion to the single-user setting [7], [8] to obtain the second of the element-wise inequalities in (68). It remains to prove the first, i.e. to show that no \( \Theta\left(\frac{1}{\sqrt{n}}\right) \) addition to \( R_{1,n} \) is possible for \( \varepsilon \in (0, 1) \). We will make use of the identities

\[
I_1(\hat{\rho}_n) = \Theta(1 - \hat{\rho}_n)
\] (71)

\[
V_1(\hat{\rho}_n) = \Theta(1 - \hat{\rho}_n)
\] (72)

which follow since \( I_1(1) = V_1(1) = 0 \), and by applying Taylor expansions in the same way as (55)–(56) (see also Appendix A). It is easily verified that the corresponding derivatives (e.g. \( \frac{\partial I_1}{\partial \rho} \)) at \( \rho = 1 \) are strictly negative, hence justifying the use of \( \Theta(\cdot) \) instead of \( O(\cdot) \) in (71)–(72).

We treat the cases \( 1 - \hat{\rho}_n = o\left(\frac{1}{n}\right) \) and \( 1 - \hat{\rho}_n = O\left(\frac{1}{n}\right) \) separately. In the former case, we combine (55)–(56) and (71)–(72) to conclude that

\[
\hat{I}_{1,n} := \mathbb{E} \left[ \frac{1}{n} \sum_{i=1}^{n} j_1(x_{1i}, x_{2i}, Y_i) \right] = \Theta(1 - \hat{\rho}_n),
\] (73)

\[
\hat{V}_{1,n} := \text{Var} \left[ \frac{1}{\sqrt{n}} \sum_{i=1}^{n} j_1(x_{1i}, x_{2i}, Y_i) \right] = \Theta(1 - \hat{\rho}_n).
\] (74)

Furthermore, we show in Appendix A that, upon replacing the given \((x_1, x_2)\) pair with a different pair yielding the same statistics of \( \sum_{i=1}^{n} j_1(x_{1i}, x_{2i}, Y_i) \) if necessary (cf. Lemma 9), we have (cf. Lemma 11)

\[
\hat{T}_{1,n} := \sum_{i=1}^{n} \mathbb{E} \left[ \left( \frac{1}{\sqrt{n}} (j_1(x_{1i}, x_{2i}, Y_i) - \mathbb{E}[j_1(x_{1i}, x_{2i}, Y_i)]) \right)^3 \right] = O\left(\frac{1 - \hat{\rho}_n}{\sqrt{n}}\right).
\] (75)

Thus, an application of the (univariate) Berry-Esseen theorem [28, Sec. XVI.5] yields

\[
\Pr \left( \frac{1}{n} \sum_{i=1}^{n} \left( j_1(x_{1i}, x_{2i}, Y_i) - \mathbb{E}[j_1(x_{1i}, x_{2i}, Y_i)] \right) \leq R_{1,n} - I_1(\hat{\rho}_n) - \gamma - \frac{\xi_1}{n} \right)
\]

\[
\geq \Phi \left( \frac{\sqrt{n}(R_{1,n} - I_1(\hat{\rho}_n) - \gamma - \frac{\xi_1}{n})}{\sqrt{\hat{V}_{1,n}} \sqrt[\sqrt{3/2}]{\hat{T}_{1,n}}} \right) - 6 \hat{T}_{1,n} \sqrt[\sqrt{3/2}]{\hat{T}_{1,n}}.
\] (76)
We see from (74)–(75) that the remainder term above scales as $O\left(\left[n(1 - \hat{\rho}_n)\right]^{-1/2}\right)$. Since it is assumed that $1 - \hat{\rho}_n = o\left(\frac{1}{n}\right)$, this remainder term vanishes. Thus, combining (76) with (70) and the fact that $\gamma = O\left(\frac{n}{\log n}\right)$, and inverting the relationship between rate and error probability, we obtain

$$R_{1,n} \leq I_1(\hat{\rho}_n) + \frac{\sqrt{\hat{V}_{1,n}}}{n} \Phi^{-1}(\varepsilon) + o\left(\frac{1}{\sqrt{n}}\right).$$

(77)

Furthermore, $\hat{\rho}_n \to 1$ implies that $\hat{V}_{1,n} \to 0$, and hence we have for any $\varepsilon \in (0,1)$ that

$$R_{1,n} \leq I_1(\hat{\rho}_n) + o\left(\frac{1}{\sqrt{n}}\right),$$

(78)

thus yielding (68) and hence (66).

It remains to show that (78) also holds when $1 - \hat{\rho}_n = O\left(\frac{1}{n}\right)$. In this case, we can combine (66) and (72) to conclude that the variance of $\sum_{i=1}^{n} j_1(x_{1i}, x_{2i}, Y_i)$ is $O(1)$, and we thus have from Chebyshev’s inequality that

$$\Pr\left(\frac{1}{n} \sum_{i=1}^{n} \left( j_1(x_{1i}, x_{2i}, Y_i) - \mathbb{E}[j_1(x_{1i}, x_{2i}, Y_i)] \right) \leq \frac{c}{\sqrt{n}} \right) \to 1$$

(79)

for all $c > 0$. Substituting (79) into (70) and taking $c \to 0$ yields (78), as desired.

6) Completion of the Proof: Combining (65) and (66), we conclude that for any sequence of codes with error probability not exceeding $\varepsilon \in (0,1)$, we have for some sequence $\hat{\rho}_n \in [-1,1]$ that

$$\left[ R_{1,n} \left[ R_{1,n} + R_{2,n} \right] \right] \in I(\hat{\rho}_n) + \frac{\Psi^{-1}(V(\hat{\rho}_n), \varepsilon)}{\sqrt{n}} + \mathcal{P}(\hat{\rho}_n, \varepsilon, n) 1,$$

(80)

where $\mathcal{P}(\rho, \varepsilon, n)$ satisfies the conditions in the theorem statement. Specifically, the first condition follows from (65) (with $\bar{g}(\rho, \varepsilon, n) := h(\rho, \varepsilon) \log \frac{n}{\varepsilon}$), and the second from (66) (with $\bar{g}(\rho, \varepsilon, n) := o\left(\frac{1}{\sqrt{n}}\right)$). This concludes the proof of the global converse.

B. Direct Part

We now prove the inner bound in (34). At a high level, we modify the key ideas in the analysis of the cost-constrained ensemble by Scarlett et al. [29] so that they are applicable to superposition codes. This approach can be applied to input-constrained MACs with DMS in significantly greater generality than the Gaussian case.

1) Random-Coding Ensemble: For user $j = 1, 2$, we introduce $K$ auxiliary cost functions $\{a_k(x_1, x_2)\}_{k=1}^{K}$, which are assumed to be arbitrary for now. The ensemble will be defined in such a way that, with probability one, each codeword pair falls into the set

$$\mathcal{D}_n := \left\{ (x_1, x_2) : \|x_1\|^2 \leq nS_1, \|x_2\|^2 \leq nS_2, \left| \frac{1}{n} \sum_{i=1}^{n} a_k(x_{1i}, x_{2i}) - \mathbb{E}[a_k(X_1, X_2)] \right| \leq \frac{\delta}{n}, \forall k \in [K] \right\},$$

(81)

where $\delta$ is a positive constant, and $(X_1, X_2)$ are jointly Gaussian with a covariance matrix of the form given in (20). Roughly speaking, the set $\mathcal{D}_n$ contains codewords satisfying the power constraints such that the empirical expectations of the $K$ auxiliary cost functions are $\delta/n$-close to the true expectations.

Before defining the ensemble, we present the following straightforward variation of [29] Prop. 1. We make use of the fact that $\|x\|^2 = \sum_{i=1}^{n} x_i^2$, i.e. the power constraints are additive in the same way as the auxiliary costs.

**Proposition 7.** Fix $\rho \in [0,1]$, and let $(X'_1, X'_2)$ be jointly distributed according to the $n$-fold product distribution of $P_{X_1, X_2} \sim \mathcal{N}(0, \Sigma(\rho))$ (see (20)), i.e. $(X'_1, X'_2) \sim \prod_{i=1}^{n} P_{X_1, X_2}(x'_{1i}, x'_{2i})$. If $\mathbb{E}[a_k(X_1, X_2)^2] < \infty$ for all $k \in [K]$, then there exists a choice of $\delta > 0$ such that

$$\Pr{( (X'_1, X'_2) \in \mathcal{D}_n )} \geq \psi(n),$$

(82)

where $\psi(n) = \Omega(n^{-(K+2)/2})$.

**Proof:** In the case that the power constraints are absent from $\mathcal{D}_n$, this is a special case of the statement of [29] Prop. 1, which was proved using the following steps: (i) Find a subset of $K' \leq K$ linearly independent auxiliary
cost functions, linear combinations of which can be used to construct the remaining $K - K'$ functions; (ii) Apply a local limit theorem (e.g. [30]) to bound the probability that all $K$ constraints in $D_n$ are satisfied. The second step relies on the constraints being two-sided, i.e. allowing for deviations on both sides of the mean in $D_n$. This is false in the presence of the power constraints. However, for $\rho \in [0, 1]$ we have $\det(\Sigma(\rho)) > 0$, thus ensuring that the functions $x_1^2$ and $x_2^2$ can be included in the set of $K' \leq K + 2$ linearly independent functions, and ensuring the validity of the second step above. For the remaining case $\rho = 1$, the codewords $X'_1$ and $X'_2$ are scalar multiples of each other, and thus the power constraints for the two users are equivalent (i.e. one implies the other). We can thus remove one of the power constraints from $D_n$ without changing the statement of the proposition, and the remaining constraint can be included in the set of $K' \leq K + 1$ linearly independent functions. This concludes the proof. ■

We now define the random-coding ensemble. We use superposition coding, in which the codewords are generated according to

$$
\left\{ \left( X_2(m_2), \{ X_1(m_1, m_2) \}_{m_1=1}^{M_1,n} \right) \right\}_{m_2=1}^{M_2,n} \sim \prod_{m_2=1}^{M_2,n} \left( P_{X_2}(x_2(m_2)) \prod_{m_1=1}^{M_1,n} P_{X_1|x_2}(x_1(m_1, m_2)|x_2(m_2)) \right) \tag{83}
$$

for some codeword distributions $P_{X_2}$ and $P_{X_1|x_2}$. These distributions will be chosen such that all codewords fall into the set $D_n$ with probability one. The auxiliary costs will be chosen to satisfy the assumptions of Proposition 7, and we assume that $\delta$ is chosen such that (82) holds, in accordance with the proposition statement.

Let the i.i.d. codewords $(X'_1, X'_2)$ be defined as in Proposition 7. Defining the set

$$
D_{2,n} := \left\{ x_2 : \Pr \left( (X'_1, X'_2) \in D_n \mid X'_2 = x_2 \right) \geq \frac{1}{2} \psi(n) \right\}, \tag{84}
$$

the codeword distributions are given by

$$
P_{X_2}(x_2) = \frac{1}{\mu_{2,n}} \prod_{i=1}^{n} P_{X_2}(x_{2i}) \mathbb{1}\left\{ x_2 \in D_{2,n} \right\}, \tag{85}
$$

$$
P_{X_1|x_2}(x_1|x_2) = \frac{1}{\mu_{1,n}(x_2)} \prod_{i=1}^{n} P_{X_1|x_2}(x_{1i}|x_{2i}) \mathbb{1}\left\{ (x_1, x_2) \in D_n \right\}, \tag{86}
$$

where $\mu_{2,n}$ and $\mu_{1,n}(x_2)$ are normalizing constants. We see that $X'_2 \in D_{2,n}$ with probability one, and the definition of $D_{2,n}$ in (84) implies that $\mu_{1,n}(x_2) \geq \frac{1}{2} \psi(n) = \Omega(n^{-(K+2)/2})$ for all $x_2 \in D_{2,n}$. It will prove useful to show that we similarly have $\mu_{2,n} = \Omega(n^{-(K+2)/2})$. To see this, we write (82) as

$$
\mu_{2,n} \Pr \left( (X'_1, X'_2) \in D_n \mid X'_2 \in D_{2,n} \right) + (1 - \mu_{2,n}) \Pr \left( (X'_1, X'_2) \in D_n \mid X'_2 \notin D_{2,n} \right) \geq \psi(n). \tag{87}
$$

Upper bounding the first probability by one and the second according to (84), we obtain

$$
\mu_{2,n} + \frac{1}{2} \psi(n) \geq \psi(n), \tag{88}
$$

where we have also used $1 - \mu_{2,n} \leq 1$. It follows that $\mu_{2,n} \geq \frac{1}{2} \psi(n) = \Omega(n^{-(K+2)/2})$, as desired. Finally, upper bounding the indicator functions in (85)–(86) by one, we obtain

$$
P_{X_1,x_2}(x_1, x_2) \leq \frac{4}{\psi(n)^2} \prod_{i=1}^{n} P_{X_1|x_2}(x_{1i}, x_{2i}). \tag{89}
$$

2) A Feinstein-type Achievability Bound for the MAC with DMS: We now state a non-asymptotic achievability based on an analogous bound for the MAC [23, Lem. 3]. This bound can be considered as a dual of Proposition 4. Define

$$
P_{X_1|x_2} W^n(y|x_2) := \int_{\mathbb{R}^n} P_{X_1|x_2}(x_1|x_2) W^n(y|x_1, x_2) \, dx_1, \tag{90}
$$

$$
P_{X_1,x_2} W^n(y) := \int_{\mathbb{R}^n} \int_{\mathbb{R}^n} P_{X_1,x_2}(x_1, x_2) W^n(y|x_1, x_2) \, dx_1 \, dx_2 \tag{91}
$$

to be output distributions induced by an input distribution $P_{X_1,x_2}$ and the channel $W^n$. We have the following proposition, whose proof is omitted since it is uses standard arguments (e.g. see [14, Thm. 4]).
Proposition 8. Fix a blocklength $n \geq 1$, random vectors $(X_1, X_2)$ with joint distribution $P_{X_1, X_2}$ such that $\|X_1\|^2_2 \leq n S_1$ and $\|X_2\|^2_2 \leq n S_2$ almost surely, auxiliary output distributions $Q_{Y|X_2}$ and $Q_Y$, and a constant $\gamma > 0$. Then there exists an $(n, M_1, n, M_2, n, S_1, S_2, \varepsilon_n)$-code for which

$$
\varepsilon_n \leq \Pr(F_n \cup G_n) + \Lambda_1 \exp(-n\gamma) + \Lambda_12 \exp(-n\gamma),
$$

where

$$
\Lambda_1 := \sup_{x_2, y} \frac{d(P_{X_1|x_2} W^n(y|x_2))}{dQ_{Y|x_2}(y|x_2)}, \quad \Lambda_12 := \sup_y \frac{d(P_{X_1, x_2} W^n(y))}{dQ_Y(y)},
$$

and

$$
F_n := \left\{ \frac{1}{n} \log \frac{W^n(Y|X_1, X_2)}{Q_{Y|x_2}(Y|x_2)} \leq \frac{1}{n} \log M_1 + \gamma \right\},$$

$$
G_n := \left\{ \frac{1}{n} \log \frac{W^n(Y|X_1, X_2)}{Q_Y(Y)} \leq \frac{1}{n} \log (M_1 M_2, n) + \gamma \right\}
$$

with $Y | \{X_1 = x_1, X_2 = x_2\} \sim W^n(\cdot|x_1, x_2)$.

The main difference between (92) and traditional Feinstein-type threshold decoding bounds (e.g. [23, Lem. 3], [31, Lem. 1]) is that we have the freedom to choose arbitrary output distributions $Q_{Y|X_2}$ and $Q_Y$; this comes at the cost of introducing the multiplicative factors $\Lambda_1$ and $\Lambda_12$ which depend on the maximum value of the Radon-Nikodym derivatives in (93). These multiplicative factors result from a standard change of measure argument.

3) Analysis of the Random-Coding Error Probability for fixed $p \in [0, 1)$: We now use Proposition 8 with the codeword input distribution $P_{X_1, X_2}$ in (85)–(86). By construction, the probability of either codeword violating the power constraint is zero. We choose the output distributions $Q_{Y|x_2} := (P_{X_1|x_2} W^n)$ and $Q_Y := (P_{X_1, x_2} W^n)$ to be of the convenient product form. As such, by (84), (86) and (89), we can choose $\Lambda_1 = 2\psi(n)^{-1}$ and $\Lambda_12 = 4\psi(n)^{-2}$ in (93). Hence,

$$
\varepsilon_n \leq 1 + \Pr\left(\frac{1}{n} \sum_{i=1}^{n} j(X_{1i}, X_{2i}, Y_i) > R_n + \gamma 1 + \frac{2 \exp(-n\gamma)}{\psi(n)} + \frac{4 \exp(-n\gamma)}{\psi(n)^2},
$$

where the information density vector $j(x_1, x_2, y)$ is defined with respect to $P_{X_1|x_2} W(y|x_2)$ and $P_{X_1, x_2} W(y)$, which coincide with $Q_{Y|x_2}$ and $Q_Y$ in (25)–(26). Choosing

$$
\gamma := \max \left\{ -\frac{1}{n} \log \left( \frac{\psi(n)}{\sqrt{n}} \right), -\frac{1}{n} \log \left( \frac{\psi(n)^2}{\sqrt{n}} \right) \right\},
$$

we notice that the sum of the third and fourth terms in (96) is upper bounded by $\frac{6}{\sqrt{n}}$. The fact that $\psi(n) = \Omega(n^{-(K+2)/2})$ (see Proposition 7) implies that $\gamma = O\left(\frac{\log n}{n}\right)$, and we obtain

$$
\varepsilon_n \leq \max_{(x_1, x_2) \in D_n} 1 - \Pr\left(\frac{1}{n} \sum_{i=1}^{n} j(x_{1i}, x_{2i}, Y_i) > R_n + \gamma 1 + \frac{6}{\sqrt{n}} \right).
$$

We now make use of the key idea proposed in [29], namely, choosing the auxiliary costs in terms of the moments of $j(x_1, x_2)$ in order to ensure that the Berry-Esseen theorem can be applied to (98). We define

$$
\mathbf{j}(x_1, x_2) := \mathbb{E}[j(x_1, x_2, Y)],
$$

$$
\mathbf{v}(x_1, x_2) := \text{Cov}\{j(x_1, x_2, Y)\},
$$

$$
\mathbf{t}(x_1, x_2) := \mathbb{E}\left[\|j(x_1, x_2, Y) - \mathbb{E}[j(x_1, x_2, Y)]\|^2\right],
$$

where the expectations and covariance are taken with respect to $W(\cdot|x_1, x_2)$. We set $K = 6$ and let the auxiliary costs equal the entries of the vectors and matrix in (99)–(101) (2 entries for $\mathbf{j}(x_1, x_2)$, 3 unique entries for the
symmetric matrix $v(x_1, x_2)$ and 1 entry for the scalar $t(x_1, x_2)$. The assumptions of Proposition 7 are satisfied, since all moments of a Gaussian random variable are finite. We have from (68), (70) and (79) that

$$\left\| E \left[ \frac{1}{n} \sum_{i=1}^{n} j(x_1, x_2, Y_i) \right] - I(\rho) \right\|_{\infty} \leq \frac{\delta}{n},$$

(102)

$$\left\| \text{Cov} \left[ \frac{1}{\sqrt{n}} \sum_{i=1}^{n} j(x_1, x_2, Y_i) \right] - V(\rho) \right\|_{\infty} \leq \frac{\delta}{n},$$

(103)

for all $(x_1, x_2) \in \mathcal{D}_n$, where the expectations and covariance are taken with respect to $W^n(\cdot|\mathbf{x})$. Similarly, defining $T(\rho) := E[t(X_1, X_2)]$, we have from (81) and (101) that

$$\left\| \frac{1}{n} \sum_{i=1}^{n} \left[ \left\| j(x_{1i}, x_{2i}, Y_i) - E[j(x_{1i}, x_{2i}, Y_i)] \right\|_{2}^{3} - T(\rho) \right] \right\|_{\infty} \leq \frac{\delta}{n},$$

(104)

Thus, by applying the multivariate Berry-Esseen theorem (26), (27) (see Appendix B) to (98) and performing Taylor expansions similarly to Section VI-A4 we obtain

$$\varepsilon_n \leq 1 - \Psi \left( \sqrt{n} \left( I_1(\rho) - R_{1,n} \right), \sqrt{n} \left( I_2(\rho) - (R_{1,n} + R_{2,n}) \right) \right) + \frac{\zeta(\rho, \delta) \log n}{\sqrt{n}},$$

(105)

where $\zeta(\rho, \delta)$ is a function depending only on $\rho$ and $\delta$. By inverting the relationship between the rates and the error probability similarly to Section VI-A4 we obtain the desired result for any given $\rho \in [0, 1)$, i.e. the first part of the theorem.

4) Analysis of the Random-Coding Error Probability for $\rho_n \to 1$: In order to prove the second part of the theorem, we use the cost-constrained ensemble with $\rho$ varying with $n$, namely $\rho_n \to 1$. Similarly to (68), it suffices to show the achievability of $(R_{1,n}, R_{2,n})$ satisfying

$$\begin{bmatrix} R_{1,n} \\ R_{1,n} + R_{2,n} \end{bmatrix} \geq I(\rho_n) + \frac{V(1)}{n} \begin{bmatrix} 0 \\ \Phi^{-1}(\varepsilon) \end{bmatrix} + o \left( \frac{1}{\sqrt{n}} \right) \mathbf{1},$$

(106)

rather than the equivalent form given by (80). See the outer bound in Lemma 6.

We set $K = 7$ and choose the first five auxiliary costs to be the same as those above, while letting $a_6$ and $a_7$ equal the two entries of

$$t(x_1, x_2) := E \left[ \left\| j(x_1, x_2, Y) - E[j(x_1, x_2, Y)] \right\|_{2}^{3} \right],$$

(107)

where the absolute value is applied element-wise. The reasons for the different choices of auxiliary costs here (compared to the case of fixed $\rho \in [0, 1)$) are to obtain a sharper bound on the third absolute moment corresponding to $\sum_{i=1}^{n} j_1(x_{1i}, x_{2i}, Y_i)$, and since we will use two applications of the scalar Berry-Esseen theorem instead of one application of the vector version.

Defining

$$T(\rho) = \begin{bmatrix} T_1(\rho) \\ T_2(\rho) \end{bmatrix} := E[t(X_1, X_2)],$$

(108)

we have similarly to (102)–(104) that

$$\left\| \frac{1}{n} \sum_{i=1}^{n} \left[ j(x_{1i}, x_{2i}, Y_i) - E[j(x_{1i}, x_{2i}, Y_i)] \right] - T(\rho) \right\|_{\infty} \leq \frac{\delta}{n}.$$

(109)

An alternative approach would be to set $K = 3$ and choose the system costs $a_1(x_1, x_2) = x_1^2, a_2(x_1, x_2) = x_2^2$ and $a_3(x_1, x_2) = x_1 x_2$. Under these choices, all codeword pairs have roughly the same powers and empirical correlations, thus allowing us to bound the moments associated with $\sum_{i=1}^{n} j(x_{1i}, x_{2i}, Y_i)$ similarly to Section VI-A. Furthermore, the uniform bound on the third absolute moment given in Lemma 10 in Appendix A can be used in place of the auxiliary cost corresponding to $t(x_1, x_2)$. On the other hand, the approach of choosing the auxiliary costs according to (99)–(101) remains applicable in significantly greater generality beyond the Gaussian setting. See [29] for more details.
We have \( I_1(1) = V_1(1) = T_1(1) = 0 \), and the behaviors of \( I_1(\rho_n) \) and \( V_1(\rho_n) \) as \( \rho_n \to 1 \) are characterized by \((71)-(72)\). Furthermore, Lemma \[12\] in Appendix \[A\] states that
\[
T_1(\rho_n) = O(1 - \rho_n) \tag{110}
\]
and hence the right-hand side of \((113)\) vanishes for any \( \rho_n \in [0,1] \), except with \( O(\cdot) \) in place of \( \Theta(\cdot) \).

Since Proposition \[7\] holds for all \( \rho \in [0,1] \) and \( \gamma \in \mathbb{R} \), it also holds for \( \rho \) varying within this range, and \((98)\) remains true with \( \gamma = O(\frac{\log n}{n}) \). Applying the union bound to \((98)\), we obtain
\[
\epsilon_n \leq \Pr \left( \frac{1}{n} \sum_{i=1}^{n} j_1(x_{1i}, x_{2i}; Y_i) \leq R_{1,n} + \gamma \right) + \Pr \left( \frac{1}{n} \sum_{i=1}^{n} j_1(x_{1i}, x_{2i}; Y_i) \leq R_{1,n} + R_{2,n} + \gamma \right) + \frac{6}{\sqrt{n}}. \tag{111}
\]
for some \( (x_1, x_2) \in \mathcal{D}_n \). The subsequent arguments are similar to Section \[VI-A5\], so we only provide an outline. We treat the cases \( 1 - \rho_n = \omega \left( \frac{1}{n} \right) \) and \( 1 - \rho_n = O \left( \frac{1}{n} \right) \) separately. For the former, we fix a small \( c > 0 \) and choose
\[
R_{1,n} = I_1(\rho_n) - \frac{c}{\sqrt{n}} - \gamma. \tag{112}
\]

Using the bounds in \((102)-(103)\) and \((109)\), and applying the (scalar) Berry-Esseen theorem, we have similarly to \((76)\) that
\[
\Pr \left( \frac{1}{n} \sum_{i=1}^{n} j_1(x_{1i}, x_{2i}; Y_i) \leq R_{1,n} + \gamma \right) \leq \Phi \left( \frac{-c}{\sqrt{n}} + O \left( \frac{\log n}{n} \right) \right) + o(1), \tag{113}
\]
where \( V_1(\rho_n) = \omega \left( \frac{1}{n} \right) \), and the remainder term is \( o(1) \) due to \((72)\) and \((110)\). Since \( \rho_n \to 1 \), we see that \( V_1(\rho_n) \to 0 \), and hence the right-hand side of \((113)\) vanishes for any \( c > 0 \). Thus, applying the Berry-Esseen theorem to \((111)\) and inverting the relationship between the rates and error probability, we obtain the following bound on the sum rate:
\[
R_{1,n} + R_{2,n} \geq I_{12}(\rho_n) + \frac{V_{12}(1)}{n} \Phi^{-1}(\epsilon) + o \left( \frac{1}{\sqrt{n}} \right). \tag{114}
\]
Combining \((112)\) and \((114)\), we obtain \((106)\) upon taking \( c \to 0 \).

Using \((102)-(103)\) and applying Chebyshev’s inequality similarly to \((79)\), we see that the left-hand side of \((113)\) also vanishes for any \( c > 0 \) in the case that \( 1 - \rho_n = O \left( \frac{1}{n} \right) \). It follows that \((106)\) remains true for this case, thus completing the proof of the second part of Theorem \[2\].

**VII. PROOF OF THEOREM 3: LOCAL SECOND-ORDER RESULT**

**A. Converse Part**

We now present the proof of the converse part of Theorem \[3\].

1) **Proof for case (i) \( (\rho = 0) \):** To prove the converse part for case (i), it suffices to consider the most optimistic case, namely \( M_{2,n} = 1 \) (i.e. no information is sent by the uninformed user). From the single-user dispersion result given in \[4, 8\] (cf. \[4\]), the number of messages for user 1 must satisfy
\[
\log M_{1,n} \leq n I_1(0) + \sqrt{n V_1(0) \Phi^{-1}(\epsilon)} + o(\sqrt{n}), \tag{115}
\]
thus proving the converse part of \((41)\).

2) **Establishing The Convergence of \( \rho_n \) to \( \rho \):** Fix a correlation coefficient \( \rho \in (0,1] \), and consider any sequence of \( (n, M_{1,n}, M_{2,n}, S_1, S_2, \varepsilon_n) \)-codes for the Gaussian MAC with DMS satisfying \((15)\). Let us consider the associated rates \( \{ (R_{1,n}, R_{2,n}) \}_{n \geq 1} \), where \( R_{j,n} = \frac{1}{n} \log M_{j,n} \) for \( j = 1, 2 \). As required by Definition \[4\], we suppose that these codes satisfy
\[
\liminf_{n \to \infty} R_{j,n} \geq R_j^*, \tag{116}
\]
\[
\liminf_{n \to \infty} \sqrt{n} (R_{j,n} - R_j^*) \geq L_j, \quad j = 1, 2, \tag{117}
\]
\[
\limsup_{n \to \infty} \varepsilon_n \leq \varepsilon \tag{118}
\]
for some \((R_1^*, R_2^*)\) on the boundary parametrized by \(\rho\), i.e. \(R_1^* = I_1(\rho)\) and \(R_1^* + R_2^* = I_{12}(\rho)\). The first-order optimality condition in \((116)\) is not explicitly required by Definition 4 but it can easily be seen that \((117)\), which is required by Definition 4, implies \((116)\). Letting \(R_n := [R_{1,n}, R_{1,n} + R_{2,n}]^T\), we have from the global converse bound in \((34)\) that there exists a sequence \(\{\rho_n\}_{n \geq 1} \subset [-1, 1]\) such that

\[
R_n \in I(\rho_n) + \frac{\Psi^{-1}(V(\rho_n), \varepsilon)}{\sqrt{n}} + \Theta(\rho_n, \varepsilon, n) 1.
\]  

(119)

Although \(\Theta(\rho_n, \varepsilon, n)\) depends on \(\rho_n\), we know from Theorem 2 that it is \(o\left(\frac{1}{\sqrt{n}}\right)\) for both \(\rho_n \to \pm 1\) and \(\rho_n\) bounded away from \(\pm 1\). It follows that

\[
R_n \in I(\rho_n) + \frac{\Psi^{-1}(V(\rho_n), \varepsilon)}{\sqrt{n}} + o\left(\frac{1}{\sqrt{n}}\right) 1.
\]  

(120)

for all sequences \(\{\rho_n\}_{n \geq 1}\). We claim that this result implies that \(\rho_n\) converges to \(\rho\). Indeed, since the boundary of the capacity region is curved and uniquely parametrized by \(\rho\) for \(\rho \in (0, 1]\), \(\rho_n \not\to \rho\) implies for some \(\delta > 0\) and for all sufficiently large \(n\) that either \(I_1(\rho_n) \leq I_1(\rho) - \delta\) or \(I_{12}(\rho_n) \leq I_{12}(\rho) - \delta\). It follows from \((120)\) that \(R_{1,n} \leq I_1(\rho_n) + \frac{\delta}{2}\) and \(R_{1,n} + R_{2,n} \leq I_{12}(\rho_n) + \frac{\delta}{2}\) for \(n\) large enough. As such, we deduce that \(R_{1,n} \leq I_1(\rho) - \frac{\delta}{2}\) or \(R_{1,n} + R_{2,n} \leq I_{12}(\rho) - \frac{\delta}{2}\) for all sufficiently large \(n\). This, in turn, contradicts the first-order optimality conditions in \((116)\).

3) Establishing The Convergence Rate of \(\rho_n\) to \(\rho\): Because each entry of \(I(\rho)\) is twice continuously differentiable, a Taylor expansion yields

\[
I(\rho_n) = I(\rho) + D(\rho)(\rho_n - \rho) + O((\rho_n - \rho)^2) 1,
\]  

(121)

where \(D(\rho)\) is the derivative of \(I\) defined in \((37)\). In the same way, since each entry of \(V(\rho)\) is continuously differentiable in \(\rho\), we have

\[
||V(\rho_n) - V(\rho)||_{\infty} = O(\rho_n - \rho).
\]  

(122)

We claim that these expansions, along with \((120)\), imply that

\[
R_n \in I(\rho) + D(\rho)(\rho_n - \rho) + \frac{\Psi^{-1}(V(\rho_n), \varepsilon)}{\sqrt{n}} + o\left(\frac{1}{\sqrt{n}}\right) + O((\rho_n - \rho)^2) + O\left(\frac{(\rho_n - \rho)^{1/2}}{\sqrt{n}}\right) 1.
\]  

(123)

The final term in the square parentheses results from the outer bound in Lemma 6 for the case \(\rho = 1\). For \(\rho \in (0, 1)\), a standard Taylor expansion yields \((123)\) with the last term replaced by \(O\left(\frac{\rho_n - \rho}{\sqrt{n}}\right)\), and it follows that \((123)\) holds for all \(\rho \in (0, 1]\).

We treat two cases separately: (I) \(\rho_n - \rho = O\left(\frac{1}{\sqrt{n}}\right)\), and (II) \(\rho_n - \rho = \omega\left(\frac{1}{\sqrt{n}}\right)\). We first show that Case (II) is not of interest in our study of local second-order coding rates (Definition 4). If Case (II) holds, intuitively in \((123), \frac{1}{\sqrt{n}}\Psi^{-1}(V(\rho), \varepsilon)\) is dominated by \(D(\rho)(\rho_n - \rho)\) and hence, the second-order term scales as \(\omega\left(\frac{1}{\sqrt{n}}\right)\) instead of the desired \(\Theta\left(\frac{1}{\sqrt{n}}\right)\). To be more precise, because

\[
\Psi^{-1}(V(\rho), \varepsilon) \subset \left[\frac{\sqrt{V_1(\rho)}\Phi^{-1}(\varepsilon)}{\sqrt{V_{12}(\rho)}\Phi^{-1}(\varepsilon)}\right]^{-},
\]  

(124)

the bound in \((123)\) implies that \(R_n\) must satisfy

\[
R_n \leq I(\rho) + D(\rho)(\rho_n - \rho) + o(\rho_n - \rho) 1.
\]  

(125)

In other words, \((123)\) and \(\rho_n - \rho = \omega\left(\frac{1}{\sqrt{n}}\right)\) imply that \(R_n\) must satisfy

\[
R_n \leq I(\rho) + D(\rho)(\rho_n - \rho) + o(\rho_n - \rho) 1.
\]  

(126)

Since the first entry of \(D(\rho)\) is negative and the second entry is positive, \((126)\) states that \(L_1 = +\infty\) (i.e. a large addition to \(R_1^*\)) only if \(L_1 + L_2 = -\infty\) (i.e. a large backoff from \(R_1^* + R_2^*\)), and \(L_1 + L_2 = +\infty\) only if \(L_1 = -\infty\). This is due the fact that we consider second-order terms scaling as \(\Theta\left(\frac{1}{\sqrt{n}}\right)\). Thus, only Case (I) is of interest, i.e. \(\rho_n - \rho = O\left(\frac{1}{\sqrt{n}}\right)\).
4) Completing the Proof for Case (ii) ($\rho \in (0, 1)$): Assuming now that $\rho_n - \rho = O\left(\frac{1}{\sqrt{n}}\right)$, it follows that $\tau_n := \sqrt{n} (\rho_n - \rho)$ is a bounded sequence. By the Bolzano-Weierstrass theorem [20 Thm. 3.6(b)], $\{\tau_n\}_{n \geq 1}$ contains a convergent subsequence $\{\tau_{n_k}\}_{k \geq 1}$. Fix any convergent subsequence $\{\tau_{n_k}\}_{k \geq 1}$ and let the limit of this subsequence be $\beta \in \mathbb{R}$, i.e., $\beta := \lim_{k \to \infty} \tau_{n_k}$. Then, for the blocklengths indexed by $n_k$, we know from (123) that

$$\sqrt{n_k} (R_{n_k} - I(\rho)) \in \beta D(\rho) + \Psi^{-1}(V(\rho), \varepsilon) + o(1) \mathbf{1},$$

where the $o(1)$ term combines the $o\left(\frac{1}{\sqrt{n}}\right)$ term in (123) and the deviation $\max\{-D_1(\rho), D_12(\rho)\}$. By referring to the second-order optimality condition in (117), and applying the definition of the limit inferior in Def. 3.16, we know that every convergent subsequence of $\{R_{j,n}\}_{n \geq 1}$ has a subsequential limit that satisfies $\lim_{k \to \infty} \sqrt{n_k} (R_{j,n_k} - R^*_j) \geq L_j$ for $j = 1, 2$. In other words, for all $\gamma > 0$, there exist an integer $K_1$ such that

$$\sqrt{n_k} (R_{1,n_k} - I_1(\rho)) \geq L_1 - \gamma \quad \text{and} \quad \sqrt{n_k} (R_{1,n_k} + R_{1,n_k} - I_{12}(\rho)) \geq L_1 + L_2 - 2\gamma$$

for all $k \geq K_1$. Thus, we may lower bound each component in the vector on the left of (127) with $L_1 - \gamma$ and $L_1 + L_2 - 2\gamma$. There also exists an integer $K_2$ such that the $o(1)$ terms are upper bounded by $\gamma$ for all $k \geq K_2$. We conclude that any pair of $(\varepsilon, R^*_1, R^*_2)$-achievable second-order coding rates $(L_1, L_2)$ must satisfy

$$\left[\frac{L_1 - 2\gamma}{L_1 + L_2 - 3\gamma}\right] \in \bigcup_{\beta \in \mathbb{R}} \left\{\beta D(\rho) + \Psi^{-1}(V(\rho), \varepsilon)\right\}.$$  

Finally, since $\gamma > 0$ is arbitrary, we can take $\gamma \downarrow 0$, thus completing the converse proof for case (ii).

5) Completing the Proof for Case (iii) ($\rho = 1$): The case $\rho = 1$ is handled in essentially the same way as $\rho \in (0, 1)$, so we only state the differences. Since $\beta$ represents the difference between $\rho_n$ and $\rho$, and since $\rho_n \leq 1$, we should only consider the case that $\beta \leq 0$. Furthermore, for $\rho = 1$ the set $\Psi^{-1}(V(\rho), \varepsilon)$ can be written in a simpler form; see Lemma 6 Using this form, we readily obtain (43).

B. Direct Part

We obtain the local result from the global result using a similar (yet simpler) argument to the converse part in Section VII-A. For fixed $\rho \in [0, 1]$ and $\beta \in \mathbb{R}$, let

$$\rho_n := \rho + \frac{\beta}{\sqrt{n}},$$

where we require $\beta \geq 0$ (resp. $\beta \leq 0$) when $\rho = 0$ (resp. $\rho = 1$). By Theorem 2 we can achieve $(R_{1,n}, R_{2,n})$ satisfying

$$R_{n} \in I(\rho_n) + \frac{\Psi^{-1}(V(\rho_n), \varepsilon)}{\sqrt{n}} + o\left(\frac{1}{\sqrt{n}}\right) \mathbf{1}.$$  

Substituting (131) into (132) and performing Taylor expansions in an identical fashion to the converse part (cf. the argument from (121) to (123)), we obtain

$$R_{n} \in I(\rho) + \frac{\beta D(\rho)}{\sqrt{n}} + \frac{\Psi^{-1}(V(\rho), \varepsilon)}{\sqrt{n}} + o\left(\frac{1}{\sqrt{n}}\right) \mathbf{1}.$$  

We immediately obtain the desired result for case (ii) where $\rho \in [0, 1)$. We also obtain the desired result for case (iii) where $\rho = 1$ using the alternative form of $\Psi^{-1}(V(1), \varepsilon)$ (see Lemma 6), similarly to the converse proof.

For case (i), we substitute $\rho = 0$ into (38) and (39) to obtain $D(\rho) = [0 \quad D_{12}(\rho)]^T$ with $D_{12}(\rho) > 0$. Since $\beta$ can be arbitrarily large, it follows from (133) that $L_2$ can take any real value. Furthermore, the set $\Psi^{-1}(V(0), \varepsilon)$ contains vectors with a first entry arbitrarily close to $\sqrt{V_1(0)}\Phi^{-1}(\varepsilon)$ (provided that the other entry is sufficiently negative), and we thus obtain (41).
APPENDIX A

MOMENTS OF THE INFORMATION DENSITY VECTOR

Let $\rho \in [-1, 1]$ be given, and recall the definition of the information density vector in (27), and the choices of $Q_{Y|X_2}$ and $Q_Y$ in (25)–(26). For a given pair of sequences $(x_1, x_2)$, form the random vector

$$A_n := \frac{1}{\sqrt{n}} \sum_{i=1}^{n} j(x_{1i}, x_{2i}, Y_i), \quad (A.1)$$

where $Y_i\{X_1 = x_{1i}, X_2 = x_{2i}\} \sim W(\cdot|x_{1i}, x_{2i})$. Define the constants $\alpha := S_1(1 - \rho^2)$, $\vartheta := S_1 + S_2 + 2\rho\sqrt{S_1S_2}$ and $\kappa := \rho \sqrt{S_1/S_2}$. Then, it can be verified that

$$j_1(x_1, x_2, Y) = \frac{1}{2} \log(1 + \alpha) - \frac{Z^2}{2} + \frac{(x_1 - \kappa x_2 + Z)^2}{2(1 + \alpha)} - \frac{\alpha Z^2 + 2(x_1 - \kappa x_2)Z}{2(1 + \alpha)} + f_1(x_1, x_2), \quad (A.2)$$

$$j_{12}(x_1, x_2, Y) = \frac{1}{2} \log(1 + \vartheta) - \frac{Z^2}{2} + \frac{(x_1 + x_2 + Z)^2}{2(1 + \vartheta)} - \frac{\vartheta Z^2 + 2(x_1 + x_2)Z}{2(1 + \vartheta)} + f_{12}(x_1, x_2), \quad (A.3)$$

where $Z := Y - x_1 - x_2 \sim \mathcal{N}(0, 1)$ and $f_1(x_1, x_2)$ and $f_{12}(x_1, x_2)$ are some deterministic functions that will not affect the covariance matrix. Taking the expectation, we obtain

$$\mathbb{E}[j_1(x_1, x_2, Y)] = \frac{1}{2} \log(1 + \alpha) - \frac{1}{2} \frac{1 + (x_1 - \kappa x_2)^2}{2(1 + \alpha)} = \frac{1}{2} \log(1 + \alpha) + \frac{(x_1 - \kappa x_2)^2 - \alpha}{2(1 + \alpha)}, \quad (A.4)$$

$$\mathbb{E}[j_{12}(x_1, x_2, Y)] = \frac{1}{2} \log(1 + \vartheta) - \frac{1}{2} \frac{1 + (x_1 + x_2)^2}{2(1 + \vartheta)} = \frac{1}{2} \log(1 + \vartheta) + \frac{(x_1 + x_2)^2 - \vartheta}{2(1 + \vartheta)}. \quad (A.5)$$

Setting $x_1 \leftarrow x_{1i}$, $x_2 \leftarrow x_{2i}$ and $Y \leftarrow Y_i$ in (A.4) and (A.5) and summing over all $i$, we conclude that the mean vector of $A_n$ is

$$\mathbb{E}[A_n] = \sqrt{n} \left[ C(\alpha) + \frac{\|x_1 - \kappa x_2\|^2 - n\alpha}{2n(1 + \alpha)} C(\vartheta) + \frac{\|x_1 + x_2\|^2 - n\vartheta}{2n(1 + \vartheta)} \right]^T. \quad (A.6)$$

From (A.2) and (A.3), we deduce that the variances are

$$\text{Var}[j_1(x_1, x_2, Y)] = \text{Var} \left[ -\alpha Z^2 + 2(x_1 - \kappa x_2)Z \right] = \frac{\alpha^2 + 2(x_1 - \kappa x_2)^2}{2(1 + \alpha)}, \quad (A.7)$$

$$\text{Var}[j_{12}(x_1, x_2, Y)] = \text{Var} \left[ -\vartheta Z^2 + 2(x_1 + x_2)Z \right] = \frac{\vartheta^2 + 2(x_1 + x_2)^2}{2(1 + \vartheta)}, \quad (A.8)$$

where we used the fact that $\text{Cov}[Z^2, Z] = \mathbb{E}Z^3 - (\mathbb{E}Z)(\mathbb{E}Z^2) = 0$. The covariance is

$$\text{Cov}[j_1(x_1, x_2, Y), j_{12}(x_1, x_2, Y)] = \text{Cov} \left[ -\alpha Z^2 + 2(x_1 - \kappa x_2)Z, -\vartheta Z^2 + 2(x_1 + x_2)Z \right] \quad (A.9)$$

$$= \frac{1}{4(1 + \alpha)(1 + \vartheta)} \left\{ \mathbb{E}((-\alpha Z^2 + 2(x_1 - \kappa x_2)Z)(-\vartheta Z^2 + 2(x_1 + x_2)Z)] \right. \quad (A.10)$$

$$\left. - \mathbb{E}(-\alpha Z^2 + 2(x_1 - \kappa x_2)Z)\mathbb{E}(-\vartheta Z^2 + 2(x_1 + x_2)Z) \right\}$$

$$= \frac{3\alpha \vartheta + 4(x_1 - \kappa x_2)(x_1 + x_2) - \alpha \vartheta}{4(1 + \alpha)(1 + \vartheta)} = \frac{\alpha \vartheta + 2(x_1^2 + (1 - \kappa)x_1x_2 - \kappa x_2^2)}{2(1 + \alpha)(1 + \vartheta)}, \quad (A.11)$$

Setting $x_1 \leftarrow x_{1i}$, $x_2 \leftarrow x_{2i}$ and $Y \leftarrow Y_i$ in (A.7), (A.8) and (A.11) and summing over all $i$, we conclude that covariance matrix of $A_n$ is

$$\text{Cov}[A_n] = \begin{bmatrix}
\frac{n\alpha^2 + 2\|x_1 - \kappa x_2\|^2}{2n(1 + \alpha)^2} & \frac{\alpha \vartheta + 2(\|x_1\|^2 + (1 - \kappa)(x_1, x_2) - \kappa \|x_2\|^2)}{2n(1 + \alpha)(1 + \vartheta)} \\
\frac{n\alpha \vartheta + 2(\|x_1\|^2 + (1 - \kappa)(x_1, x_2) - \kappa \|x_2\|^2)}{2n(1 + \alpha)(1 + \vartheta)} & \frac{\vartheta \alpha^2 + 2\|x_1 + x_2\|^2}{2n(1 + \vartheta)^2}
\end{bmatrix}. \quad (A.12)
In the remainder of the section, we analyze several third absolute moments associated with $A_n$ appearing in the univariate [28 Sec. XVI.5] and multivariate Berry-Esseen theorems [26, 27] (see Appendix B). The following lemma will be used to replace any given (x₁, x₂) pair by an “equivalent” pair (in the sense that the statistics of $A_n$ are unchanged) for which the corresponding third moments have the desired behavior. This is analogous to Polyanskiy et al. [8], where for the AWGN channel, one can use a spherical symmetry argument to replace any given sequence $x$ such that $\|x\|_2^2 = nS$ with a fixed sequence $(\sqrt{S}, \ldots, \sqrt{S})$.

**Lemma 9.** The joint distribution of $A_n$ depends on $(x_1, x_2)$ only through the powers $\|x_1\|_2^2$, $\|x_2\|_2^2$ and the inner product $(x_1, x_2)$.

**Proof:** This follows by substituting (A.2)–(A.3) into (A.1) and using the symmetry of the additive noise sequence $Z = (Z_1, \ldots, Z_n)$. For example, from (A.2), the first entry of $A_n$ can be written as

$$\frac{1}{\sqrt{n}} \left( \frac{1}{2} \log(1 + \alpha) - \frac{1}{2} \|Z\|_2^2 + \frac{1}{2(1 + \alpha)}\|x_1 - \kappa x_2 + Z\|_2^2 \right),$$

and the desired result follows by writing

$$\|x_1 - \kappa x_2 + Z\|_2^2 = \|x_1\|_2^2 + \kappa^2 \|x_2\|_2^2 + \|Z\|_2^2 - 2\kappa(x_1, x_2) + 2(x_1, Z) - 2\kappa(x_2, Z).$$

Since $Z$ is i.i.d. Gaussian, the last two terms depend on $(x_1, x_2)$ only through $\|x_1\|_2^2$ and $\|x_2\|_2^2$.

We now provide lemmas showing that, upon replacing a given pair $(x_1, x_2)$ with an equivalent pair using Lemma 9 if necessary, the corresponding third moments have the desired behavior. It will prove useful to work with the empirical correlation coefficient

$$\rho_{\text{emp}}(x_1, x_2) := \frac{\langle x_1, x_2 \rangle}{\|x_1\|_2 \|x_2\|_2}.$$  

(A.15)

It is easily seen that Lemma 9 remains true when the inner product $(x_1, x_2)$ is replaced by this normalized quantity.

**Lemma 10.** For any fixed $\tilde{\rho} \in [-1, 1]$, there exists a sequence of pairs $(x_1, x_2)$ (indexed by increasing lengths $n$) such that $\|x_1\|_2^2 = nS_1$, $\|x_2\|_2^2 = nS_2$, $\rho_{\text{emp}}(x_1, x_2) = \tilde{\rho}$, and

$$\tilde{T}_n := \sum_{i=1}^{n} \mathbb{E} \left[ \left\| \frac{1}{\sqrt{n}} (j(x_{1i}, x_{2i}, Y_i) - \mathbb{E}[j(x_{1i}, x_{2i}, Y_i)]) \right\|_2^3 \right] = O \left( \frac{1}{\sqrt{n}} \right),$$

(A.16)

where the $O(\frac{1}{\sqrt{n}})$ term is uniform in $\tilde{\rho} \in [-1, 1]$.

**Proof:** By using the fact that $\|v\|_2 \leq \|v\|_1$ and $(|a| + |b|)^3 \leq 4|a|^3 + 4|b|^3$, the following bounds hold:

$$\tilde{T}_n \leq \sum_{i=1}^{n} \mathbb{E} \left[ \left\| \frac{1}{\sqrt{n}} (j(x_{1i}, x_{2i}, Y_i) - \mathbb{E}[j(x_{1i}, x_{2i}, Y_i)]) \right\|_1^3 \right]$$

(A.17)

$$\leq 4 \sum_{i=1}^{n} \mathbb{E} \left[ \left\| \frac{1}{\sqrt{n}} (j_1(x_{1i}, x_{2i}, Y_i) - \mathbb{E}[j_1(x_{1i}, x_{2i}, Y_i)]) \right\|_1^3 \right]$$

$$+ 4 \sum_{i=1}^{n} \mathbb{E} \left[ \left\| \frac{1}{\sqrt{n}} (j_{12}(x_{1i}, x_{2i}, Y_i) - \mathbb{E}[j_{12}(x_{1i}, x_{2i}, Y_i)]) \right\|_1^3 \right].$$

(A.18)

We now specify $(x_1, x_2)$ whose powers and correlation match those given in the lemma statement. Assuming for the time being that $|\tilde{\rho}| \leq \frac{n-1}{n}$, we choose

$$x_1 = (\sqrt{S_1}, \ldots, \sqrt{S_1})$$

(A.19)

$$x_2 = (\sqrt{S_2(1 + \eta)}, \sqrt{S_2}, \ldots, \sqrt{S_2}, -\sqrt{S_2(1 - \eta)}, -\sqrt{S_2}, \ldots, -\sqrt{S_2})$$

(A.20)

where $\eta \in (-1, 1)$, and $x_2$ contains $k \geq 1$ negative entries and $n - k \geq 1$ positive entries. It is easily seen that $\|x_1\|_2^2 = nS_1$ and $\|x_2\|_2^2 = nS_2$, as desired. Furthermore, we can choose $k$ and $\eta$ to obtain the desired correlation since

$$\langle x_1, x_2 \rangle = (n - 2(k - 1) + \sqrt{1 + \eta} - \sqrt{1 - \eta}) \sqrt{S_1S_2},$$

(A.21)
and since the range of the function $f(\eta) := \sqrt{1 + \eta} - \sqrt{1 - \eta}$ for $\eta \in (-1, 1)$ is given by $(-\sqrt{2}, \sqrt{2})$.

Using (A.2)–(A.3), it can easily be verified that the third absolute moment of each entry of $j(x_1, x_2, Y)$ (i.e. $\mathbb{E}|j_1(x_1, x_2, Y) - \mathbb{E}[j_1(x_1, x_2, Y)]|^3$ and $\mathbb{E}|j_2(x_1, x_2, Y) - \mathbb{E}[j_2(x_1, x_2, Y)]|^3$) is bounded above by some constant for any $(x_1, x_2) = (\sqrt{S_1}, \pm \sqrt{S_2})$ ($c \in (0, 2)$). We thus obtain (A.16) using (A.18). The proof is concluded by noting that a similar argument applies for the case $\tilde{\rho} \in \left(\frac{n-1}{n}, 1\right]$ by replacing (A.20) by

$$x_2 = \left(\sqrt{S_2}(1 + \eta), \sqrt{S_2}(1 - \eta), \sqrt{S_2}, \cdots, \sqrt{S_2}\right),$$  

(A.22)

and similarly (with negative entries) when $\tilde{\rho} \in [-1, \frac{n-1}{n})$.

\textbf{Lemma 11.} For any pair of sequences $\{\rho_n\}_{n \geq 1} \subset [-1, 1]$ and $\{\tilde{\rho}_n\}_{n \geq 1} \subset [-1, 1]$ such that $\rho_n, \tilde{\rho}_n \to 1$ with $1 - \tilde{\rho}_n = \Theta(1 - \rho_n)$, there exists a sequence of vectors $(x_1, x_2)$ (indexed by increasing lengths $n$) with the properties $\|x_1\|^2 = nS_1$, $\|x_2\|^2 = nS_2$, $\rho_{\text{emp}}(x_1, x_2) = \tilde{\rho}_n$, and they satisfy

$$\tilde{T}_{1,n} := \sum_{i=1}^{n} \mathbb{E}\left[\left|\frac{1}{\sqrt{n}}(j_1(x_{1,i}, x_{2,i}, Y_i) - \mathbb{E}[j_1(x_{1,i}, x_{2,i}, Y_i)])\right|^3\right] = O\left(\frac{1 - \tilde{\rho}_n}{\sqrt{n}}\right),$$  

(A.23)

where the information density is computed with respect to $\rho_n$, i.e. $j_1$ is defined in (27) with input distribution $P_{X_1, X_2} = \mathcal{N}(0, \Sigma(\rho_n))$.

\textbf{Proof:} We make use of the notation and definitions in the proof of Lemma 10 using a subscript to denote dependence on $n$, e.g. $k_n$. We again focus on the case that $|\rho_n| \leq \frac{n-1}{n}$ and we choose $(x_1, x_2)$ according to (A.19)–(A.20). The remaining cases are handled similarly (e.g. using (A.22)) and are thus omitted.

Inverting (A.21), we have $k_n = \frac{\alpha}{2}(1 - \rho_n) + O(1)$, and thus the entries corresponding to $x_2 \neq \sqrt{S_2}$ in (A.20) contribute an additive $O(\frac{1 - \rho_n}{\sqrt{n}})$ term to the summation in (A.23). For the remaining entries, where $(x_1, x_2) = (\sqrt{S_1}, \sqrt{S_2})$, we note from (A.2) that the third absolute moment of $j(x_1, x_2, Y)$ is given by

$$\mathbb{E}\left[|j_1(x_1, x_2, Y) - \mathbb{E}[j_1(x_1, x_2, Y)]|^3\right] = \mathbb{E}\left[\frac{-\alpha(Z^2 - 1) + 2(x_1 - \kappa x_2)Z^2}{2(1 + \alpha)}\right]^3 \leq 2\left(|\alpha|^3\mathbb{E}[Z^2 - 1]^3\right) + 8\left(\sqrt{S_1} - \kappa \sqrt{S_2}\right)^3\mathbb{E}[Z]^3 \leq \Theta\left(|\alpha|^3 + (\sqrt{S_1} - \kappa \sqrt{S_2})^3\right) \leq \Theta(1 - \tilde{\rho}_n),$$  

(A.24)

(A.25)

(A.26)

where (A.25) follows because $\alpha \geq 0$ and $|a+b|^3 \leq 4(|a|^3 + |b|^3)$, and (A.27) follows by substituting $\alpha = S_1(1 - \rho_n)$ and $\kappa = \rho \sqrt{S_1/S_2}$, and applying first-order Taylor expansions in $1 - \tilde{\rho}_n$ (recall that $1 - \rho_n = \Theta(1 - \tilde{\rho}_n)$ by assumption). From (A.27), we conclude that the entries where $(x_1, x_2) = (\sqrt{S_1}, \sqrt{S_2})$ in (A.23) contribute an additive $O(\frac{1 - \rho_n}{\sqrt{n}})$ term to the summation in (A.23), thus yielding the desired result.

Recall the definition of $T_1(\rho)$ in (108) which we restate here for the reader’s convenience:

$$T_1(\rho) := \mathbb{E}\left[\mathbb{E}\left[|j_1(X_1, X_2, Y) - \mathbb{E}[j_1(X_1, X_2, Y)]|^3 | X_1, X_2\right]\right].$$  

(A.28)

The distribution of $(X_1, X_2)$ above is $\mathcal{N}(0, \Sigma(\rho))$ and the information density $j_1$ is also defined with respect to $\rho$.

\textbf{Lemma 12.} For any sequence $\{\rho_n\}_{n \geq 1}$ satisfying $\rho_n \to 1$, we have

$$T_1(\rho_n) = O(1 - \rho_n).$$  

(A.29)

\textbf{Proof:} We upper bound $t_1(x_1, x_2)$ using the Cauchy-Schwarz and arithmetic-geometric mean inequalities as follows:

$$t_1(x_1, x_2) = \mathbb{E}\left[|j_1(x_1, x_2, Y) - \mathbb{E}[j_1(x_1, x_2, Y)]|^3\right] \leq \sqrt{\mathbb{E}\left[|j_1(x_1, x_2, Y) - \mathbb{E}[j_1(x_1, x_2, Y)]|^2\right]} \mathbb{E}\left[(j_1(x_1, x_2, Y) - \mathbb{E}[j_1(x_1, x_2, Y)])^4\right] \leq \frac{1}{2} \left(\mathbb{E}\left[|j_1(x_1, x_2, Y) - \mathbb{E}[j_1(x_1, x_2, Y)]|^2\right] + \mathbb{E}\left[(j_1(x_1, x_2, Y) - \mathbb{E}[j_1(x_1, x_2, Y)])^4\right]\right) \cdot (A.30)$$

(A.31)

(A.32)
Now, by \((A.7)\), the variance of \(j_1(x_1, x_2, Y)\) is
\[
m_2(x_1, x_2) := \mathbb{E} \left[ (j_1(x_1, x_2, Y) - \mathbb{E}[j_1(x_1, x_2, Y)])^2 \right] = \frac{\alpha^2 + 2(x_1 - \kappa x_2)^2}{2(1 + \alpha)^2} \tag{A.33}
\]
By a similar calculation, we deduce that the centralized fourth moment of \(j_1(x_1, x_2, Y)\) is
\[
m_4(x_1, x_2) := \mathbb{E} \left[ (j_1(x_1, x_2, Y) - \mathbb{E}[j_1(x_1, x_2, Y)])^4 \right] = \frac{15\alpha^4 + 60\alpha^2(x_1 - \kappa x_2)^2 + 12(x_1 - \kappa x_2)^4}{4(1 + \alpha)^4} \tag{A.34}
\]
Taking the expectation of \((A.33)\) with respect to \((X_1, X_2) \sim \mathcal{N}(0, \Sigma(\rho))\) yields
\[
\mathbb{E}[m_2(X_1, X_2)] = \frac{\alpha(2 + \alpha)}{2(1 + \alpha)^2} \tag{A.35}
\]
where we used the fact that \(X_1 - \kappa X_2 \sim \mathcal{N}(0, \alpha)\). Similarly, taking the expectation of \((A.34)\) yields
\[
\mathbb{E}[m_4(X_1, X_2)] = \frac{3\alpha^2(5\alpha^2 + 20\alpha + 12)}{4(1 + \alpha)^4} \tag{A.36}
\]
where we used the fact that \(\mathbb{E}[(X_1 - \kappa X_2)^4] = 3\alpha^2\). Now, observe that since \(\alpha = \alpha(\rho) = S_1(1 - \rho^2)\) is continuously differentiable in \(\rho\), so are \(\mathbb{E}[m_2(X_1, X_2)]\) and \(\mathbb{E}[m_4(X_1, X_2)]\). Moreover \(\mathbb{E}[m_2(X_1, X_2)]\) and \(\mathbb{E}[m_4(X_1, X_2)]\) are equal to zero when \(\rho = 1\) (and hence \(\alpha = 0\)). Consequently, by \((A.32)\) we know that \(T_1(\rho)\) is upper bounded by a continuously differentiable function in \(\rho\) and it evaluates to zero when \(\rho = 1\). We conclude, by first-order Taylor expansions in \(\rho\) in the vicinity of \(\rho = 1\), that \((A.29)\) holds.

**APPENDIX B**

**THE MULTIVARIATE BERRY-ESSEEN THEOREM**

In this section, we state a version of the multivariate Berry-Esseen theorem \([26], [27]\) that is suited for our needs in this paper. The following is a restatement of Corollary 38 in \([32]\).

**Theorem 13.** Let \(U_1, \ldots, U_n\) be independent, zero-mean random vectors in \(\mathbb{R}^d\). Let \(L_n := \frac{1}{\sqrt{n}}(U_1 + \cdots + U_n)\). Assume \(V := \text{Cov}(L_n)\) is positive definite with minimum eigenvalue \(\lambda_{\text{min}}(V) > 0\). Let \(t := \frac{1}{n} \sum_{i=1}^{n} \mathbb{E} \|U_i\|^2\) and let \(Z\) be a zero-mean Gaussian random vector with covariance \(V\). Then, for all \(n \in \mathbb{N}\),
\[
\sup_{\mathcal{C} \in \mathcal{C}_d} \left| \Pr(L_n \in \mathcal{C}) - \Pr(Z \in \mathcal{C}) \right| \leq \frac{k_d t}{\lambda_{\text{min}}(V)^{3/2} \sqrt{n}} \tag{B.1}
\]
where \(\mathcal{C}_d\) is the family of all convex, Borel measurable subsets of \(\mathbb{R}^d\), and \(k_d\) is a function only of the dimension \(d\) (e.g., \(k_2 = 265\)).

**APPENDIX C**

**PROOF OF LEMMA 5**

Fix \((z_1, z_2) \in \Psi^{-1}(V, \varepsilon + \lambda_n)\) and define \(Z = (Z_1, Z_2) \sim \mathcal{N}(0, V)\). Since \(\Psi^{-1}(V, \varepsilon)\) is monotonic in the sense that \(\Psi^{-1}(V, \varepsilon) \subset \Psi^{-1}(V, \varepsilon')\) for \(\varepsilon \leq \varepsilon'\), it suffices to verify that \((z_1, z_2)\) belongs to the set on the right-hand side of \((63)\) for those \((z_1, z_2)\) on the boundary of \(\Psi^{-1}(V, \varepsilon + \lambda_n)\). That is (cf. \((31)\)),
\[
\Pr \left( Z_1 \leq -z_1, Z_2 \leq -z_2 \right) = 1 - (\varepsilon + \lambda_n) \tag{C.1}
\]
Define \(\nu_n := \inf \{ \nu > 0 : (-z_1 - \nu, -z_2 - \nu) \in \Psi^{-1}(V, \varepsilon) \}\). We need to show that \(\nu_n = o(1)\) is bounded above by some linear function of \(\lambda_n\). By using \((C.1)\) and the definition of \(\nu_n\), we see that
\[
\lambda_n = \Pr \left( Z_1 \in [-z_1 - \nu_n, -z_1] \cup Z_2 \in [-z_2 - \nu_n, -z_2] \right) \tag{C.2}
\]
\[
\geq \max_{j=1,2} \left\{ \Phi \left( \frac{-z_j}{\sqrt{V_{jj}}} \right) - \Phi \left( \frac{-z_j - \nu_n}{\sqrt{V_{jj}}} \right) \right\} \tag{C.3}
\]
The assumption that $V$ is a non-zero positive-semidefinite matrix ensures that at least one of $V_{jj}, j = 1, 2$ is non-zero. We have the lower bound
\[
\Phi \left( \frac{-z}{\sqrt{V}} \right) - \Phi \left( \frac{-z - \nu_n}{\sqrt{V}} \right) \geq \frac{\nu_n}{\sqrt{V}} \min \{ \mathcal{N}(z; 0, V), \mathcal{N}(z + \nu_n; 0, V) \}.
\] (C.4)

Hence, for all $n$ large enough, each of the terms in $\{ \}$ in (C.3) is bounded below by $\nu_n f(z_j, V_{jj})$ for $j = 1, 2$ where $f(z, V) := \frac{1}{2\sqrt{V}} \mathcal{N}(z; 0, V)$ satisfies $\lim_{z \to \pm \infty} f(z, V) = 0$. Hence, $\nu_n \leq \lambda_n \min_{j=1,2} f(z_j, V_{jj}^{-1})$. For every fixed $\varepsilon \in (0,1)$, every $(z_1, z_2) \in \Psi^{-1}(V, \varepsilon + \lambda_n)$ satisfies $\min \{|z_1|, |z_2| \} < \infty$, and hence $\min_{j=1,2} f(z_j, V_{jj}^{-1})$ is finite. This concludes the proof.

**Appendix D**

**Proof of Lemma 6**

Recall that $\hat{\rho}_n \to 1$. We start by proving the inner bound on $\Psi^{-1}(V(\hat{\rho}_n), \varepsilon)$. Let $(w_1, w_2)$ be an arbitrary element of the left-hand-side of (67), i.e. $w_1 \leq -b_n$ and $w_2 \leq \sqrt{V_{12}(1)} \Phi^{-1}(\varepsilon + \lambda_n) - b_n$. Define the random variables $(Z_{1,n}, Z_{2,n}) \sim \mathcal{N}(0, V(\hat{\rho}_n))$ and the sequence $b_n := (1 - \hat{\rho}_n)^{1/4}$. Consider
\[
\Pr (Z_{1,n} \leq -w_1, Z_{2,n} \leq -w_2) \geq \Pr (Z_{1,n} \leq b_n, Z_{2,n} \leq -\left( \sqrt{V_{12}(1)} \Phi^{-1}(\varepsilon + \lambda_n) - b_n \right)) \tag{D.1}
\]
\[
\geq \Pr (Z_{2,n} \leq -\left( \sqrt{V_{12}(1)} \Phi^{-1}(\varepsilon + \lambda_n) - b_n \right)) - \Pr (Z_{1,n} > b_n) \tag{D.2}
\]
\[
= \Phi \left( -\left( \sqrt{V_{12}(1)} \Phi^{-1}(\varepsilon + \lambda_n) - b_n \right) \right) - \Phi \left( \frac{-b_n}{\sqrt{V_{12}(\hat{\rho}_n)}} \right). \tag{D.3}
\]

From the choice of $b_n$ and the fact that $\sqrt{V_{1}(\hat{\rho}_n)} = \Theta(\sqrt{1 - \hat{\rho}_n})$ (cf. (72)), the argument of the second term scales as $-(1 - \hat{\rho}_n)^{-1/4}$ which tends to $-\infty$. Hence, the second term vanishes. We may thus choose a vanishing sequence $a_n$ so that the expression in (D.3) equals $1 - \varepsilon$. Such a choice satisfies $a_n = \Theta(b_n) = \Theta((1 - \hat{\rho}_n)^{1/4})$, in accordance with the lemma statement. From the definition in (31), we have proved that $(w_1, w_2) \in \Psi^{-1}(V(\hat{\rho}_n), \varepsilon)$ for this choice of $(a_n, b_n)$.

For the outer bound on $\Psi^{-1}(V(\hat{\rho}_n), \varepsilon)$, let $(u_1, u_2)$ be an arbitrary element of $\Psi^{-1}(V(\hat{\rho}_n), \varepsilon)$. By definition,
\[
\Pr (Z_{1,n} \leq -u_1, Z_{2,n} \leq -u_2) \geq 1 - \varepsilon, \tag{D.4}
\]
where $(Z_{1,n}, Z_{2,n}) \sim \mathcal{N}(0, V(\hat{\rho}_n))$ as above. Thus,
\[
1 - \varepsilon \leq \Pr (Z_{2,n} \leq -u_2) = \Phi \left( \frac{-u_2}{\sqrt{V_{12}(\hat{\rho}_n)}} \right). \tag{D.5}
\]
This leads to
\[
u_n \leq \sqrt{V_{12}(\hat{\rho}_n)} \Phi^{-1}(\varepsilon) = \sqrt{V_{12}(1)} \Phi^{-1}(\varepsilon) + c_n' \tag{D.6}
\]
for some $c_n' = \Theta(1 - \hat{\rho}_n)$, since $\rho \mapsto \sqrt{V_{12}(\rho)}$ is continuously differentiable and its derivative does not vanish at $\rho = 1$. Similarly, we have
\[
u_n \leq \sqrt{V_{1}(\hat{\rho}_n)} \Phi^{-1}(\varepsilon) = c_n'' \tag{D.7}
\]
for some $c_n'' = \Theta(\sqrt{1 - \hat{\rho}_n})$, since $V_1(1) = 0$ and $\sqrt{V_1(\hat{\rho}_n)} = \Theta(\sqrt{1 - \hat{\rho}_n})$ by (72). Letting $c_n := \max \{|c_n'|, |c_n''|\} = \Theta(\sqrt{1 - \hat{\rho}_n})$, we deduce that $(u_1, u_2)$ belongs to the rightmost set in (67). This completes the proof.

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