Two Reflected Gray Code based orders
on some restricted growth sequences

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Abstract
We consider two order relations: that induced by the $m$-ary reflected Gray code and a
suffix partitioned variation of it. We show that both of them when applied to some sets of
restricted growth sequences still yield Gray codes. These sets of sequences are: subexcedant
or ascent sequences, restricted growth functions, and staircase words. In each case we give
efficient exhaustive generating algorithms and compare the obtained results.

1 Introduction and motivations
The term ‘Gray code’ was taken from Frank Gray, who patented Binary Reflected Gray Code
(BRGC) in 1953 [1]. The concept of BRGC is extended to Reflected Gray Code (RGC), to
accommodate $m$-tuples (sequence), with $m > 2$ [2]. In these Gray codes, successive sequence
differ in a single position, and by $+1$ or $-1$ in this position. More generally, if a list of sequences
is such that the Hamming distance between successive sequences is upper bounded by a constant
$d$, then the list is said a $d$-Gray code. So in particular, BRGC and RGC are 1-Gray codes. In
addition, if the positions where the successive sequences differ are adjacent, then we say that
the list is a $d$-adjacent Gray code.

For long time, the design of Gray codes for combinatorial classes and their corresponding
generating algorithms was an ad-hoc task, that is, done case by case according to the class
under consideration. Recently, general techniques which fit to large classes of combinatorial
objects were developed and used. Among them are, for example, the ECO-method (initiated in
Gray code context in [3], see also [1]), prefix rotations (yielding bubble languages, see [5][6]
and references therein), or Reflected Gray Code based order relations; this last technique was used
implicitly, for example in [7][8], and developed systematically as a general method in [4][9][13].
The results presented in this paper are in the light of this last direction. More precisely, we show
that two order relations induced by Reflected Gray Code and its variation also give Gray codes
for some classes of restricted growth sequences defined by means of statistics. These classes
are: subexcedant or ascent sequences, restricted growth functions, and staircase words. We give
efficient (CAT) generating algorithm for each obtained Gray code.
2 Preliminaries

2.1 Gray code orders

Let $G_n(m)$ be the set of length $n$ $m$-ary sequences $s_1 s_2 \ldots s_n$ with $s_i \in \{0, 1, \ldots, m - 1\}$; clearly, $G_n(m)$ is the product set $\{0, 1, \ldots, m - 1\}^n$. The Reflected Gray Code (RGC for short) for the set $G_n(m)$, denoted by $G_n(m)$, is the natural extension of the Binary Reflected Gray Code to this set. The list $G_n(m)$ is defined recursively by the following relation \[2]:

$$G_n(m) = \begin{cases} 
    \epsilon & \text{if } n = 0, \\
    \mathcal{G}_{n-1}(m), 1\mathcal{G}_{n-1}(m), 2\mathcal{G}_{n-1}(m), \ldots, (m-1)\mathcal{G}_{n-1}(m) & \text{if } n > 0,
\end{cases}$$

where $\epsilon$ is the empty sequence, $\mathcal{G}_{n-1}(m)$ is the reverse of $G_{n-1}(m)$, and $\mathcal{G}'_{n-1}(m)$ is $G_{n-1}(m)$ or $\mathcal{G}_{n-1}(m)$ according to $m$ is odd or even.

In $G_n(m)$, two successive sequences differ in a single position and by $+1$ or $-1$ in this position. A list for a set of sequences induces an order relation to this set, and we give two order relations induced by the RGC and its variation, namely RGC order \[10\] and Co-RGC order.

We adopt the convention that lower case bold letters represent tuples, for example: $s = s_1 s_2 \ldots s_n$, $a = a_1 a_2 \ldots a_k$, $b = b_{k+1} b_{k+2} \ldots b_n$.

**Definition 1.** The Reflected Gray Code order $\prec$ on $G_n(m)$ is defined as: $s = s_1 s_2 \ldots s_n$ is less than $t = t_1 t_2 \ldots t_n$, denoted by $s \prec t$, if either

- $\sum_{i=1}^{k-1} s_i$ is even and $s_k < t_k$, or
- $\sum_{i=1}^{k-1} s_i$ is odd and $s_k > t_k$,

where $k$ is the leftmost position where $s$ and $t$ differ.

It is easy to see that $G_n(m)$ defined in relation \[10\] lists sequences in $G_n(m)$ in $\prec$ order.

Now we give a variation of $G_n(m)$. Let $s_1 s_2 \ldots s_n$ be a sequence in $G_n(m)$. The complement of $s_i$, $1 \leq i \leq n$, is

$$(m-1-s_i),$$

and the reverse of $s_1 s_2 \ldots s_n$ is

$$s_n s_{n-1} \ldots s_1.$$

Let $\tilde{G}_n(m)$ be the list obtained by transforming each sequence $s$ in $G_n(m)$ as follows:

- complementing each digit in $s$ if $m$ is even, or complementing only digits in odd positions if $m$ is odd, then
- reversing the obtained sequence.

Clearly, $\tilde{G}_n(m)$ is also a Gray code for $G_n(m)$, and the sequences therein are listed in Co-Reflected Gray Code order, as defined formally below.

**Definition 2.** The Co-Reflected Gray Code order $\prec_c$ on $G_n(m)$ is defined as: $s = s_1 s_2 \ldots s_n$ is less than $t = t_1 t_2 \ldots t_n$, denoted by $s \prec_c t$, if either

- $\sum_{i=k+1}^{n} s_i + (n-k)$ is even and $s_k > t_k$, or

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• \[ \sum_{i=k+1}^{n} s_i + (n - k) \] is odd and \( s_k < t_k \),

where \( k \) is the rightmost position where \( s \) and \( t \) differ.

Although this definition sounds somewhat arbitrary, as we will see in Section 4, it turns out that \( \prec_c \) order gives suffix partitioned Gray codes for some sets of restricted growth sequences. Obviously, the restriction of \( G_n(m) \) (resp. \( \tilde{G}_n(m) \)) to a set of sequences is simply the list of sequences in the set listed in \( \prec \) (resp. \( \prec_c \)) order.

2.2 Restricted growth sequences defined by means of statistics

Through this paper we consider sequences over non-negative integers. A statistic on a set of sequences is an association of an integer to each sequence in the set. For a sequence \( s_1 s_2 \ldots s_n \), its length minus one, numbers of ascents/levels/descents, maximal value, and last value are classical examples of statistics. They are defined as follows, see also [14]:

- \( \text{len}(s_1 s_2 \ldots s_n) = n - 1; \)
- \( \text{asc}(s_1 s_2 \ldots s_n) = \text{card}\{i \mid 1 \leq i < n \text{ and } s_i < s_{i+1}\}; \)
- \( \text{lev}(s_1 s_2 \ldots s_n) = \text{card}\{i \mid 1 \leq i < n \text{ and } s_i = s_{i+1}\}; \)
- \( \text{des}(s_1 s_2 \ldots s_n) = \text{card}\{i \mid 1 \leq i < n \text{ and } s_i > s_{i+1}\}; \)
- \( m(s_1 s_2 \ldots s_n) = \max\{s_1, s_2, \ldots, s_n\}; \)
- \( \text{lv}(s_1 s_2 \ldots s_n) = s_n. \)

If \( st \) is one of the statistics \( \text{len}, \text{asc}, m, \) and \( \text{lv} \), then \( st \) satisfy the following:

\[ st(s_1 s_2 \ldots s_n) \leq n - 1, \] (2)

and

\[ \text{if } s_n = st(s_1 s_2 \ldots s_{n-1}) + 1, \text{ then } s_n = st(s_1 s_2 \ldots s_{n-1}s_n). \] (3)

On the contrary, the statistics \( \text{lev} \) and \( \text{des} \) do not satisfy relation (3). Accordingly, through this paper we will consider only the four statistics above. However, as we will point out, some of the results presented here are also true for arbitrary statistics satisfying relations (2) and (3).

Definition 3. For a given statistic \( st \), an \( st \)-restricted growth sequence \( s_1 s_2 \ldots s_n \) is a sequence with \( s_1 = 0 \) and

\[ 0 \leq s_{k+1} \leq st(s_1 s_2 \ldots s_k) + 1 \text{ for } 1 \leq k < n, \] (4)

and the set of \( st \)-restricted growth sequences is the set of all sequences \( s_1 s_2 \ldots s_n \) satisfying relation (4).

From this definition, it follows that any prefix of an \( st \)-restricted growth sequence is also (a shorter) \( st \)-restricted growth sequence.

Remark 1. If \( st \) is a statistic satisfying relations (2) and (3) above, then

1. \( \max\{st(s_1 s_2 \ldots s_n) \mid s_1 s_2 \ldots s_n \text{ is an } st \text{-restricted growth sequence}\} = n - 1; \)
2. if \( s_1 s_2 \ldots s_n \) is an \( st \)-restricted growth sequence, then for any \( k, 1 \leq k < n \), \( s_{k+1} = \text{st}(s_1 s_2 \ldots s_k) + 1 \) implies \( s_{k+1} = \text{st}(s_1 s_2 \ldots s_k s_{k+1}) \).

The sets of \( st \)-restricted growth sequences, where \( st \) is one of the statistics \( \text{len}, \text{asc}, m, \) and \( lv \), are defined below.

**Definition 4.**

- The set \( SE_n \) of subexcedant sequences of length \( n \) is defined as:
  \[
  SE_n = \{ s_1 s_2 \ldots s_n \mid s_1 = 0 \text{ and } 0 \leq s_{k+1} \leq \text{len}(s_1 s_2 \ldots s_k) + 1 \text{ for } 1 \leq k < n \};
  \]
- The set \( A_n \) of ascent sequences of length \( n \) is defined as:
  \[
  A_n = \{ s_1 s_2 \ldots s_n \mid s_1 = 0 \text{ and } 0 \leq s_{k+1} \leq \text{asc}(s_1 s_2 \ldots s_k) + 1 \text{ for } 1 \leq k < n \};
  \]
- The set \( R_n \) of restricted growth functions of length \( n \) is defined as:
  \[
  R_n = \{ s_1 s_2 \ldots s_n \mid s_1 = 0 \text{ and } 0 \leq s_{k+1} \leq m(s_1 s_2 \ldots s_k) + 1 \text{ for } 1 \leq k < n \};
  \]
- The set \( S_n \) of staircase words of length \( n \) is defined as:
  \[
  S_n = \{ s_1 s_2 \ldots s_n \mid s_1 = 0 \text{ and } 0 \leq s_{k+1} \leq lv(s_1 s_2 \ldots s_k) + 1 \text{ for } 1 \leq k < n \}.
  \]

Notice that alternatively, \( SE_n = \{0\} \times \{0, 1\} \times \ldots \times \{0, 1, \ldots, n - 1\} \).

**Remark 2.** \( S_n \subset R_n \subset A_n \subset SE_n \subset G_n(n) \).

Below we give examples to illustrate Remark 2.

**Example 1.**

- If \( s = 010145 \), then \( s \in SE_6 \), but \( s \notin A_6 \), \( s \notin R_6 \) and \( s \notin S_6 \).
- If \( s = 010103 \), then \( s \in SE_6 \), \( s \in A_6 \), but \( s \notin R_6 \) and \( s \notin S_6 \).
- If \( s = 010102 \), then \( s \in SE_6 \), \( s \in A_6 \), and \( s \in R_6 \), but \( s \notin S_6 \).
- If \( s = 010101 \), then \( s \in SE_6 \), \( s \in A_6 \), \( s \in R_6 \) and \( s \in S_6 \).

See Table 3 for the sets \( S_5, R_5, \) and \( A_5 \) listed in \( \prec \) order, and Table 4 for the same sets listed in \( \prec_c \) order.

We denote by

- \( X_n \) the list for the set \( X_n \) in \( \prec \) order, and by \( \tilde{X}_n \) that in \( \prec_c \) order;
- \( \text{succ}_X(s) \), \( s \in X_n \), the successor of \( s \) in the set \( X_n \) listed in \( \prec \) order; that is, the smallest sequence in \( X_n \) larger than \( s \) with respect to \( \prec \) order;
- \( \tilde{\text{succ}}_X(s) \) the counterpart of \( \text{succ}_X(s) \) with respect to \( \prec_c \) order;
- \( \text{first}(\mathcal{L}) \) the first sequence in the list \( \mathcal{L} \);
- \( \text{last}(\mathcal{L}) \) the last sequence in the list \( \mathcal{L} \).
2.3 Constant Amortized Time algorithms and principle

An exhaustive generating algorithm is said to run in constant amortized time (CAT for short) if the total amount of computation is proportional to the number of generated objects. And so, a CAT algorithm can be considered an efficient algorithm.

Ruskey and van Baronaigien [15] introduced three CAT properties, and proved that if a recursive generating procedure satisfies them, then it runs in constant amortized time (see also [16]). They called this general technique to prove the efficiency of a generating algorithm as CAT principle, and the involved properties are:

1. Every call of the procedure results in the output of at least one object;
2. Excluding the computation done by the recursive calls, the amount of computation of any call is proportional to the degree of the call, that is, the number of call initiated by the current call;
3. The number of calls of degree one, if any, is $O(N)$, where $N$ is the number of generated objects.

All the generating algorithms we present in this paper satisfy these three desiderata, and so they are efficient.

3 The Reflected Gray Code order for the sets $SE_n$, $A_n$, $R_n$, and $S_n$

3.1 The bound of Hamming distance between successive sequences in the lists $SE_n$, $A_n$, $R_n$, and $S_n$

Here we will show that the Hamming distance between two successive sequences in each of the mentioned lists is upper bounded by a constant, and so the lists are Gray codes.

Without another specification, $X_n$ generically denotes one of the sets $SE_n$, $A_n$, $R_n$, or $S_n$; and $X_n$ denotes its corresponding list in $\prec$ order, that is, one of the lists $SE_n$, $A_n$, $R_n$, or $S_n$. Later in this section, Theorem 1 and Proposition 2 state that, in each case, the set $X_n$ listed in $\prec$ order yields a Gray code.

**Lemma 1.** If $s = s_1s_2\ldots s_n$ and $t = t_1t_2\ldots t_n$ are two sequences in $X_n$ with $t = \text{succ}_X(s)$ and $k$ is the leftmost position where they differ, then $s_k = t_k + 1$ or $s_k = t_k - 1$.

**Proof.** Let $t = \text{succ}_X(s)$ and $k$ be the leftmost position where they differ. Let us suppose that $s_k < t_k$ and $s_k \neq t_k - 1$ (the case $s_k > t_k$ and $s_k \neq t_k + 1$ being similar).

It is easy to check that $u = s_1s_2\ldots s_{k-1}(s_k + 1)0\ldots 0$ belongs to $X_n$, and considering the definition of $\prec$ order relation, it follows that $s \prec u \prec t$, which is in contradiction with $t = \text{succ}_X(s)$, and the statement holds.

If $a = a_1a_2\ldots a_k \in X_k$, then for any $n > k$, $a$ is the prefix of at least one sequence in $X_n$, and we denote by $a \mid X_n$ the sublist of $X_n$ of all sequences having the prefix $a$. Clearly, a list in $\prec$ order for a set of sequences is a *prefix partitioned* list (all sequences with same prefix are contiguous), and for any $a \in X_k$ and $n > k$, it follows that $a \mid X_n$ is a contiguous sublist of $X_n$.
For a given \( a \in X_k \), the set of all \( x \) such that \( ax \in X_{k+1} \) is called the defining set of the prefix \( a \), and obviously \( ax \) is also a prefix of some sequences in \( X_n \), for any \( n > k \). We denote by

\[
\omega_X(a) = \max\{x \mid ax \in X_{k+1}\}
\]

the largest value in the defining set of \( a \). And if we denote \( \omega_X(a) \) by \( M \), then by Remark \( \square \) we have

\[
M = \text{st}(a) + 1 = \text{st}(aM).
\]

And consequently,

\[
\omega_X(aM) = \text{st}(aM) + 1 = M + 1.
\]

The next proposition gives the pattern of \( s \in X_n \), if \( s = \text{last}(a \mid X_n) \) or \( s = \text{first}(a \mid X_n) \).

**Proposition 1.** Let \( k < n \) and \( a = a_1a_2\ldots a_k \in X_k \). If \( s = \text{last}(a \mid X_n) \), then the pattern of \( s \) is given by:

- if \( \sum_{i=1}^{k} a_i \) is odd, then \( s = a0\ldots0 \);
- if \( \sum_{i=1}^{k} a_i \) is even and \( M \) is odd, then \( s = aM0\ldots0 \);
- if \( \sum_{i=1}^{k} a_i \) is even and \( M \) is even, then \( s = aM(M+1)0\ldots0 \);

where \( M \) denotes \( \omega_X(a) \).

Similar results hold for \( s = \text{first}(a \mid X_n) \) by replacing ‘odd’ by ‘even’, and vice versa, for the parity of \( \sum_{i=1}^{k} a_i \).

**Proof.** Let \( s = a_1a_2\ldots a_ks_{k+1}\ldots s_n = \text{last}(a \mid X_n) \).

If \( \sum_{i=1}^{k} a_i \) is odd, then by considering the definition of \( < \) order, it follows that \( s_{k+1} \) is the smallest value in the defining set of \( a \), and so \( s_{k+1} = 0 \), and finally \( s = a0\ldots0 \), and the first point holds. Now let us suppose that \( \sum_{i=1}^{k} a_i \) is even. In this case \( s_{k+1} \) equals \( \omega_X(a) = M \), the largest value in the defining set of \( a \). When in addition \( M \) is odd, so is the summation of \( aM \), the length \( k+1 \) prefix of \( s \), and thus \( s = aM0\ldots0 \), and the second point holds.

Finally, when \( M \) is even, then \( s_{k+2} \) is the largest value in the defining set of \( aM \), which by relation \( \square \) is \( M+1 \). In this case \( M+1 \) is odd, and thus \( s = aM(M+1)0\ldots0 \), and the last point holds.

The proof for the case \( s = \text{first}(a \mid X_n) \) is similar. \( \square \)

By Proposition \( \square \) above, we have the following:

**Theorem 1.** The lists \( A_n, R_n \) and \( S_n \) are 3-adjacent Gray codes.

**Proof.** Let \( X_n \) be one of the lists \( A_n, R_n \) or \( S_n \), and \( t = \text{succ}_X(s) \). Let \( k \) be the leftmost position where \( s \) and \( t \) differ, and let us denote by \( a \) the length \( k \) prefix of \( s \) and \( a' \) that of \( t \); so, \( s = \text{last}(a \mid X_n) \) and \( t = \text{first}(a' \mid X_n) \). If \( k + 3 \leq n \), then by Proposition \( \square \) it follows that \( s_{k+3} = s_{k+4} = \cdots = s_n = 0 \) and \( t_{k+3} = t_{k+4} = \cdots = t_n = 0 \). So \( s \) and \( t \) differ only in position \( k \), and possibly in position \( k + 1 \) and in position \( k + 2 \).

Now we show the adjacency, that is, if \( k + 2 \leq n \) and \( s_{k+1} = t_{k+1} \) implies \( s_{k+2} = t_{k+2} \). If \( s_{k+1} = t_{k+1} \), by Lemma \( \square \) it follows that the summation of the length \( k \) prefix of \( s \) and that of \( t \) have different parity, and two cases can occur:
• $s_{k+1} = t_{k+1} = 0$, and by Proposition 1 it follows that $s_{k+2} = t_{k+2} = 0$; or
• $s_{k+1} = t_{k+1} \neq 0$, and thus $s_{k+1} = t_{k+1} = \omega(a) = \omega(a')$. In this case, $\omega(a)$ is either odd and so $s_{k+2} = t_{k+2} = 0$, or is even and so $s_{k+2} = t_{k+2} = \omega(a) + 1$.

In both cases, $s_{k+2} = t_{k+2}$.

It is well known that the restriction of $G_n(m)$ defined in relation (1) to any product space remains a 1-Gray code, see for example [13]. In particular, for $SE_n = \{0\} \times \{0, 1\} \times \ldots \times \{0, 1, \ldots, n-1\}$ we have the next proposition. Its proof is simply based on Lemma 1, Proposition 1 and on the additional remark: for any $a \in SE_k$, $k < n$, it follows that $\omega_{SE}(a) = k$.

**Proposition 2.** The list $SE_n$ is 1-Gray code.

It is worth to mention that for any statistic $st$ satisfying relations (2) and (3), the list in $\prec$ order for the set of $st$-restricted growth sequences of length $n$ is an at most 3-Gray code. Actually, the lists $SE_n$, $A_n$, $R_n$, and $S_n$ are circular Gray codes, that is, the last and the first sequences in the list differ in the same way. Indeed, by the definition of $\prec$ order, it follows that:

- $first(X_n) = 000\ldots 0$;
- $last(X_n) = 010\ldots 0$;

where $X_n$ is one of the list $SE_n$, $A_n$, $R_n$, or $S_n$.

### 3.2 Generating algorithms for the lists $SE_n$, $A_n$, $R_n$, and $S_n$

Procedure $Gen1$ in Figure 1 is a general procedure generating exhaustively the list of $st$-restricted growth sequences, where $st$ is a statistic satisfying relations (2) and (3). According to particular instances of the function $\Omega_X$ called by it (and so, of the statistic $st$), $Gen1$ produces specific $st$-restricted growth sequences, and in particular the lists $SE_n$, $A_n$, $R_n$, and $S_n$. From the length one sequence 0, $Gen1$ constructs recursively increasing length $st$-restricted growth sequences: for a given prefix $s_1s_2\ldots s_k$ it produces all prefixes $s_1s_2\ldots s_{k+i}$, with $i$ covering (in increasing or decreasing order) the defining set of $s_1s_2\ldots s_k$; and eventually all length $n$ $st$-restricted growth sequences. It has the following parameters:

- $k$, the position in the sequence $s$ which is updated by the current call;
- $x$, belongs to the defining set of $s_1s_2\ldots s_{k-1}$, and is the value to be assigned to $s_k$;
- $dir$, the direction (ascending for $dir \mod 2 = 0$ and descending for $dir \mod 2 = 1$) to cover the defining set of $s_1s_2\ldots s_{k-1}$;
- $v$, the value of the statistic of the prefix $s_1s_2\ldots s_{k-1}$ from which the value of the statistic of the current prefix $s_1s_2\ldots s_k$ is computed. Remark that $v = \omega_X(s_1s_2\ldots s_{k-1}) - 1$.

Function $\Omega_X$ computes $\omega_X(s_1s_2\ldots s_k)$ (see relation (5)), and the main call is $Gen1(1, 0, 0, 0)$. In $Gen1$ the amount of computation of each call is proportional with the degree of the call, and there are no degree one calls, and so it satisfies the CAT principle stated at the end of Section 2 and so it is an efficient generating algorithm. The computational tree of $Gen1$ producing the list $A_4$ is given in Figure 2. Each node at level $k$, $1 \leq k \leq 4$, represents prefixes $s_1s_2\ldots s_k$, and leaves sequences in $A_4$. 

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procedure Gen1(k, x, dir, v)
global: n, s;
    sk := x;
    if k = n
        then Print s;
    else u := Omega_X(v, k);
        if dir mod 2 = 0
            then for i := 0 to u do
                Gen1(k + 1, i, i, u - 1);
            enddo
        else for i := u downto 0 do
                Gen1(k + 1, i, i + 1, u - 1);
            enddo
        endif
    end.

(i) function Omega_SE(w, p)
    return p;
end.

(ii) function Omega_A(w, p)
    if p > 1 and sp > sp-1
        then return wp + 2;
    else return wp + 1;
    endif
end.

(iii) function Omega_R(w, p)
    if p > 1 and sp > w
        then return wp + 2;
    else return wp + 1;
    endif
end.

(iv) function Omega_S(w, p)
    return sp + 1;
end.

Figure 1: (a) Algorithm Gen1, generating the list \(X_n\); (b) Particular function Omega_X called by Gen1, and returning the value for \(\omega_X(s_1s_2\ldots s_k)\), if \(X_n\) is one of the sets: (i) \(SE_n\), (ii) \(A_n\), (iii) \(R_n\), and (iv) \(S_n\).

Figure 2: The tree induced by the initial call Gen1(1, 0, 0, 0) for \(n = 4\) and generating the list \(A_4\).
4 The Co-Reflected Gray Code order for the sets $SE_n$, $A_n$, $R_n$, and $S_n$

In this section we will consider, as in the previous one, the sets $SE_n$, $A_n$, $R_n$, and $S_n$, but listed in $\prec_c$ order. Our main goal is to prove that the obtained lists are Gray codes as well, and to develop generating algorithms for these lists. Recall that $X_n$ generically denotes one of the sets $SE_n$, $A_n$, $R_n$, or $S_n$; and let $\mathcal{X}_n$ denote their corresponding list in $\prec_c$ order, that are, $\mathcal{SE}_n$, $\mathcal{A}_n$, $\mathcal{R}_n$, or $\mathcal{S}_n$. Clearly, a set of sequences listed in $\prec_c$ order is a suffix partitioned list, that is, all sequences with same suffix are contiguous, and such are the lists we consider here.

For a set $X_n$ and a sequence $b = b_{k+1}b_{k+2} \ldots b_n$, we call $b$ an admissible suffix in $X_n$ if there exists at least a sequence in $X_n$ having suffix $b$. For example, $124$ is an admissible suffix in $A_6$, because there are sequences in $A_6$ ending with $124$, namely $012124$ and $010124$. On the other hand, $224$ is not an admissible suffix in $A_6$; indeed, there is no length $6$ ascent sequence ending with $224$.

We denote by $\mathcal{X}_n \mid b$ the sublist of $\mathcal{X}_n$ of all sequences having suffix $b$, and clearly, $\mathcal{X}_n \mid b$ is a contiguous sublist of $\mathcal{X}_n$. The set of all $x$ such that $xb$ is also an admissible suffix in $X_n$ is called the defining set of the suffix $b$.

For $\prec$ order discussed in Section 3, the characterization of prefixes is straightforward: $a_1a_2 \ldots a_k$ is the prefix of some sequences in $X_n$, $n > k$, if and only if $a_1a_2 \ldots a_k$ is in $X_k$. And the defining set of the prefix $a_1a_2 \ldots a_k$ is $\{0, 1, \ldots, \text{st}(a_1a_2 \ldots a_k) + 1\}$. In the case of $\prec_c$ order, it turns out that similar notions are more complicated: for example, $13$ is an admissible suffix in $A_5$, but $13$ is not in $A_2$; and the defining set of the suffix $13$ is $\{0, 2\}$, because $013$ and $213$ are both admissible suffixes in $A_5$, but $113$ is not. See Table II for the set $A_5$ listed in $\prec_c$ order.

4.1 Suffix expansion of sequences in the sets $SE_n$, $A_n$, $R_n$, and $S_n$

For a suffix partitioned list, we need to build st-restricted growth sequences under consideration from right to left, i.e., by expanding their suffix. For this purpose, we need the notions defined below.

**Definition 5.** Let $b = b_{k+1}b_{k+2} \ldots b_n$, $1 \leq k < n$, an admissible suffix in $X_n$.

- $\alpha_X(b)$ is the set of all elements in the defining set of the suffix $b$. Formally:

$$\alpha_X(b) = \{ x \mid xb \text{ is an admissible suffix in } X_n \},$$

and for the empty suffix $\epsilon$, $\alpha_X(\epsilon) = \{0, 1, \ldots, n - 1\}$.

- $\mu_X(b)$ is the minimum required value of the statistic defining the set $X_n$, and provided by a length $(k + 1)$ prefix of a sequence in $X_n$ having suffix $b$. Formally:

$$\mu_X(b) = \min\{ \text{st}(s_1s_2 \ldots s_kb_{k+1}) \mid s_1s_2 \ldots s_k b \in X_n \}.$$

Notice that $\mu_X(xb) \in \{ \mu_X(b) - 1, \mu_X(b), x \}$ for $x \in \alpha_X(b)$.

**Remark 3.** Let $\text{st}$ be one of statistics $\text{asc}$, $m$, or $\text{lv}$, and $s = s_1s_2 \ldots s_n$ be an $\text{st}$-restricted growth sequence. If there is a $k < n$ such that $s_{k+1} = k$, then $s_i = i - 1$ for all $i$, $1 \leq i \leq k$. 

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Proposition 4. Let \( s_k < k - 1 \), then in each case for \( st, st(s_1s_2\ldots s_k) < k - 1 \), which is in contradiction with \( s_{k+1} = k \), and so \( s_k = k - 1 \). Similarly, \( s_{k-1} = k - 2, \ldots, s_2 = 1 \), and \( s_1 = 0 \). □

Under the conditions in the previous remark, \( s_{k+1} = k \) imposes that all values at the left of \( k + 1 \) in \( s \) are uniquely determined. As we will see later, in the induced tree of the generating algorithm, all descendants of a node with \( s_{k+1} = k \) have degree one, and we will eliminate the obtained degree-one path in order not to alter the algorithm efficiency.

It is routine to check the following propositions. (Actually, Proposition 3 is a consequence of Remark 1.)

Proposition 3. Let \( X_n \) be one of the sets \( SE_n, A_n, R_n, \) or \( S_n \). If \( b = b_{k+1}b_{k+2}\ldots b_n, 1 \leq k < n \), is an admissible suffix in \( X_n \), then \( b_{k+1} \leq \mu_X(b) \).

Proposition 4. Let \( Y_n \) be one of the sets \( A_n, R_n, \) or \( S_n \). If \( b = b_{k+1}b_{k+2}\ldots b_n, 1 \leq k < n \), is an admissible suffix in \( Y_n \), then

1. if \( b = b_n \), that is, a length one admissible suffix, then \( \mu_Y(b) = b_n \);
2. \( \mu_Y(b) = k \) if and only if \( b_{k+1} = 1 \);
3. if \( x b \) is also an admissible suffix in \( Y_n \) (i.e., \( x \in \alpha_Y(b) \)) and \( x \geq b_{k+1} \), then
\[
\mu_Y(xb) = \max\{x, \mu_Y(b)\}.
\]

The following propositions give the values for \( \alpha_X(b) \) and \( \mu_X(xb) \), if \( X_n \) is one of the sets \( SE_n, A_n, R_n, \) or \( S_n \). We do not provide the proofs for Propositions 5, 6, 11, and 12 because they are obviously based on the definition of the corresponding sequences.

Proposition 5. Let \( b = \epsilon \) or \( b = b_{k+1}b_{k+2}\ldots b_n \) be an admissible suffix in \( SE_n \). Then
\[
\alpha_{SE}(b) = \begin{cases} 
{0, 1, \ldots, n - 1} & \text{if } b = \epsilon, \\
{0, 1, \ldots, k - 1} & \text{otherwise}.
\end{cases}
\]

Proposition 6. Let \( b = b_{k+1}b_{k+2}\ldots b_n \) be an admissible suffix in \( SE_n \) and \( x \in \alpha_{SE}(b) \). Then
\[
\mu_{SE}(xb) = \mu_{SE}(b) - 1.
\]

Obviously, for a length one suffix \( b = b_n \), it follows that \( \mu_{SE}(b) = n - 1 \).

Example 2. If \( b = \epsilon \), and \( n = 10 \), then \( \alpha_{SE}(b) = \{0, 1, \ldots, 9\} \);
and for \( b = b_{10}, b_{10} \in \{0, 1, \ldots, 9\} \), it follows that \( \mu_{SE}(xb) = 9 - 1 = 8 \), for all \( x \in \alpha_{SE}(b) \).

Proposition 7. Let \( b = \epsilon \) or \( b = b_{k+1}b_{k+2}\ldots b_n \) be an admissible suffix in \( A_n \). Then
\[
\alpha_A(b) = \begin{cases} 
{0, 1, \ldots, n - 1} & \text{if } b = \epsilon, \\
\{k - 1\} & \text{if } \mu_A(b) = k, \text{ or } \mu_A(b) = k - 1 \text{ and } b_{k+1} = 0, \\
\{0, 1, \ldots, b_{k+1} - 1\} \cup \{k - 1\} & \text{if } \mu_A(b) = k - 1 \text{ and } 0 < b_{k+1} < k, \\
\{0, 1, \ldots, k - 1\} & \text{if } \mu_A(b) < k - 1.
\end{cases}
\]
Proof. If \( b = \epsilon \), the result is obvious.

For \( b \neq \epsilon \), let \( x \in \alpha_A(b) \).

If \( \mu_A(b) = k \), by Proposition 4 point 2, \( b_{k+1} = k \) and by Remark 3 we have \( x = k - 1 \).

If \( \mu_A(b) = k - 1 \) and \( b_{k+1} = 0 \), then \( \text{asc}(xb_{k+1}) = 0 \), and so \( \mu_A(xb) = \mu_A(b) = k - 1 \), and again by Proposition 4 point 2 we have \( x = k - 1 \).

If \( \mu_A(b) = k - 1 \) and \( 0 < b_{k+1} < k \), then there are two possibilities for \( \mu_A(xb) \):

- \( \mu_A(xb) = \mu_A(b) = k - 1 \), if \( \text{asc}(xb_{k+1}) = 0 \), and as above \( x = k - 1 \);
- \( \mu_A(xb) = \mu_A(b) - 1 = k - 2 \), if \( \text{asc}(xb_{k+1}) = 1 \). In this case \( x \in \{0, 1, \ldots, b_{k+1} - 1\} \).

If \( \mu_A(b) < k - 1 \) (and consequently \( 0 \leq b_{k+1} < k \)), then there are two possibilities for \( \mu_A(xb) \):

- \( \mu_A(xb) = \mu_A(b) \), if \( \text{asc}(xb_{k+1}) = 0 \), and we have \( x \in \{b_{k+1}, b_{k+1} + 1, \ldots, k\} \);
- \( \mu_A(xb) = \mu_A(b) - 1 \), if \( \text{asc}(xb_{k+1}) = 1 \), and we have \( x \in \{0, 1, \ldots, b_{k+1} - 1\} \).

\( \square \)

Proposition 8. Let \( b = b_{k+1}b_{k+2} \ldots b_n \) be an admissible suffix in \( A_n \) and \( x \in \alpha_A(b) \). Then

\[
\mu_A(xb) = \begin{cases} 
  x & \text{if } x \geq \mu_A(b), \\
  \mu_A(b) & \text{if } b_{k+1} \leq x < \mu_A(b), \\
  \mu_A(b) - 1 & \text{if } x < b_{k+1}.
\end{cases}
\]

Proof. If \( x \geq \mu_A(b) \), by Proposition 3 it follows that \( x \geq b_{k+1} \), and by Proposition 4 point 3, that \( \mu_A(xb) = \max\{x, \mu_A(b)\} = x \).

If \( b_{k+1} \leq x < \mu_A(b) \), then, again by Proposition 4 point 3, it follows that \( \mu_A(xb) = \max\{x, \mu_A(b)\} = \mu_A(b) \).

If \( x < b_{k+1} \), then \( \text{asc}(xb_{k+1}) = 1 \), so \( \mu_A(xb) = \mu_A(b) - 1 \).

\( \square \)

Example 3. Let \( k = 5 \), \( n = 9 \), and \( b = b_6b_7b_8b_9 = 2050 \) be an admissible suffix in \( A_9 \). Clearly, \( \mu_A(b) \), the minimum number of ascents in a prefix \( s_1s_2 \ldots s_5b \) such that \( s_1s_2 \ldots s_5b \in A_9 \), is 4. In this case, denoting \( s_5 \) by \( x \), we have

- the set \( \alpha_A(b) \) of all possible values for \( x \) is \( \{0, 1, \ldots, b_{k+1} - 1\} \cup \{k - 1\} = \{0, 1\} \cup \{4\} \);
- \( \mu_A(xb) = \mu_A(b) - 1 = 4 - 1 = 3 \), if \( x \in \{0, 1\} \); or \( \mu_A(xb) = \mu_A(b) = 4 \), if \( x = 4 \).

Proposition 9. Let \( b = \epsilon \) or \( b = b_{k+1}b_{k+2} \ldots b_n \) be an admissible suffix in \( R_n \). Then

\[
\alpha_R(b) = \begin{cases} 
  \{0, 1, \ldots, n-1\} & \text{if } b = \epsilon, \\
  \{k - 1\} & \text{if } \mu_R(b) = k, \text{ or } \mu_R(b) = \mu_R(b) = k - 1 \text{ and } b_{k+1} < k - 1, \\
  \{0, 1, \ldots, k - 1\} & \text{if } \mu_R(b) = k - 1 \text{ and } b_{k+1} = k - 1, \text{ or } \mu_R(b) < k - 1.
\end{cases}
\]

Proof. If \( b = \epsilon \), the result is obvious.

For \( b \neq \epsilon \), let \( x \in \alpha_R(b) \).

If \( \mu_R(b) = k \), by Proposition 4 point 2, \( b_{k+1} = k \) and by Remark 3 we have \( x = k - 1 \).

If \( \mu_R(b) = k - 1 \) and \( b_{k+1} < k - 1 \), then the maximal value of the statistic \( m \) (defining the set \( R_n \)) of a length \( k + 1 \) prefix ending with \( b_{k+1} < k - 1 \) is \( k - 1 \), and it is reached when \( x = k - 1 \).

If \( \mu_R(b) = k - 1 \) and \( b_{k+1} = k - 1 \), then there are two possibilities for \( \mu_R(xb) \):

\( \square \)
• \( \mu_R(xb) = \mu_R(b) = k - 1 \), and as above, this implies \( x = k - 1 \);

• \( \mu_R(xb) = \mu_R(b) - 1 = k - 2 \), which implies \( x \in \{0, 1, \ldots, k - 2\} \).

Finally, if \( \mu_R(b) < k - 1 \), then \( x \) can be any value in \( \{0, 1, \ldots, k - 1\} \). \( \square \)

**Proposition 10.** Let \( b = b_{k+1}b_{k+2} \ldots b_n \) be an admissible suffix in \( R_n \) and \( x \in \alpha_R(b) \). Then

\[
\mu_R(xb) = \begin{cases} 
  x & \text{if } x \geq \mu_R(b), \\
  \mu_R(b) & \text{if } b_{k+1} \leq x < \mu_R(b) \text{ or } x < b_{k+1} < \mu_R(b), \\
  \mu_R(b) - 1 & \text{if } x < b_{k+1} = \mu_R(b).
\end{cases}
\]

**Proof.** The case \( x \geq \mu_R(b) \) is analogous with the similar case in Proposition 5.

The next case is equivalent with \( x < \mu_R(b) \) and \( b_{k+1} < \mu_R(b) \), and since \( R_n \) corresponds to the statistic \( m \), the result holds.

Finally, if \( x < b_{k+1} = \mu_R(b) \), then \( \mu_R(b) = \mu_R(xb) + 1 \), and so \( \mu_R(xb) = \mu_R(b) - 1 \). \( \square \)

**Example 4.** Let \( k = 4, n = 7 \), and \( b = b_5b_6b_7 = 241 \) be an admissible suffix in \( R_7 \). It follows that \( b_{k+1} = 2, \mu_R(b) = 3 \) and

- \( \alpha_R(b) = \{k - 1\} = \{3\} \);
- \( \mu_R(xb) = x = 3 \).

**Proposition 11.** Let \( b = \epsilon \) or \( b = b_{k+1}b_{k+2} \ldots b_n \) be an admissible suffix in \( S_n \). Then

\[
\alpha_S(b) = \begin{cases} 
  \{0, 1, \ldots, n - 1\} & \text{if } b = \epsilon, \\
  \{k - 1\} & \text{if } b_{k+1} = k, \\
  \{C, C + 1, \ldots, k - 1\} & \text{if } 0 \leq b_{k+1} \leq k - 1,
\end{cases}
\]

where \( C = \max\{0, b_{k+1} - 1\} \).

Since \( \mu_S(b) = b_{k+1} \), the next result follows:

**Proposition 12.** Let \( b = b_{k+1}b_{k+2} \ldots b_n \) be an admissible suffix in \( S_n \) and \( x \in \alpha_S(b) \). Then

\[ \mu_S(xb) = x. \]

**Example 5.** Let \( k = 6, n = 9 \), and \( b = b_7b_8b_9 = 457 \) be an admissible suffix in \( S_9 \). So we have \( \mu_S(b) = 4, C = \max\{0, b_{k+1} - 1\} = \max\{0, 3\} = 3, \) and

- \( \alpha_S(b) = \{C, C + 1, \ldots, k - 1\} = \{3, 4, 5\} \);
- \( \mu_S(xb) = x, \) where \( x \in \{3, 4, 5\} \).

### 4.2 The bound of Hamming distance between successive sequences in the lists \( \mathcal{SE}_n, \mathcal{A}_n, \mathcal{R}_n, \) and \( \mathcal{S}_n \)

Now we show that the Hamming distance between two successive sequences in the mentioned lists is upper bounded by a constant, which implies that the lists are Gray codes. These results are embodied in Theorems 2, 4 and 4.
4.2.1 The list $\tilde{S}E_n$

**Theorem 2.** The list $\tilde{S}E_n$ is 1-Gray code.

**Proof.** The result follows from the fact that the restriction of the 1-Gray code list $G_n(n)$ to any product space remains a 1-Gray code (see [13]), in particular to the set

$$V_n = \vartheta_1 \times \vartheta_2 \times \ldots \times \vartheta_n,$$

where

- $\vartheta_i = \{0, 1, \ldots, n - i\}$, if $n$ is odd and $i$ is even, or
- $\vartheta_i = \{i - 1, i, \ldots, n - 1\}$, if $n$ is even, or $n$ and $i$ are both odd.

Then by applying to each sequence $s$ in the list $V_n$ the two transforms mentioned before Definition 2, namely:

- complementing each digit in $s$ if $n$ is even, or only digits in odd positions if $n$ is odd, then
- reversing the obtained sequence,

the desired 1-Gray code for the set $SE_n$ in $\prec_c$ order is obtained. \square

4.2.2 The lists $\tilde{A}_n$ and $\tilde{R}_n$

The next proposition describes the pattern of $s = \text{last}(\tilde{Y}_n | b)$ and $s = \text{first}(\tilde{Y}_n | b)$, where $\tilde{Y}_n$ is one of the lists $\tilde{A}_n$ or $\tilde{R}_n$.

**Proposition 13.** Let $Y_n$ be one of the sets $A_n$ or $R_n$, and $b = b_{k+1} b_{k+2} \ldots b_n$ be an admissible suffix in $Y_n$. If $s = \text{last}(\tilde{Y}_n | b)$ or $s = \text{first}(\tilde{Y}_n | b)$, then $s$ has one of the following patterns:

- $s = 012 \ldots (k-2)(k-1)b$, or
- $s = 012 \ldots (k-2)0b$.

**Proof.** Let $s = s_1 s_2 \ldots s_k b_{k+1} b_{k+2} \ldots b_n$. Since $s = \text{last}(\tilde{Y}_n | b)$ or $s = \text{first}(\tilde{Y}_n | b)$, according to $\alpha_Y(b)$ given in Propositions 7 and 9, it follows that $s_k \in \{0, k-1\}$. In other words, $s_k$ is either the smallest or the largest value in $\alpha_Y(b)$.

If $s_k = k - 1$, then by Remark 3 we have $s = 012 \ldots (k-2)(k-1)b$.

If $s_k = 0$, then considering the definition of $\prec_c$ order we have either

- $s = \text{first}(\tilde{Y}_n | b)$ and $\sum_{i=k+1}^n b_i + (n-k)$ is odd, or
- $s = \text{last}(\tilde{Y}_n | b)$ and $\sum_{i=k+1}^n b_i + (n-k)$ is even.

For the first case, again by the definition of $\prec_c$ order, it follows that $s_{k-1}$ must be the largest value in $\alpha_Y(0b)$, and so $s_{k-1} = k - 2$, and by Remark 3, $s = 012 \ldots (k-2)0b$. Similarly, the same result is obtained for the second case. \square

A direct consequence of the previous proposition is the next theorem.

**Theorem 3.** The lists $\tilde{A}_n$ and $\tilde{R}_n$ are 2-adjacent Gray codes.
Proof. Let \( s, t \in Y_n \), with \( \tau = \text{succy}(s) \). If \( k + 1 \) is the rightmost position where \( s \) and \( t \) differ, then there are admissible suffixes \( b = b_{k+1} b_{k+2} \ldots b_n \) and \( b' = b'_{k+1} b_{k+2} \ldots b_n \) in \( Y_n \) such that \( s = \text{last}(\tilde{\gamma}_n | b) \) and \( t = \text{first}(\tilde{\gamma}_n | b') \).

By Proposition 13, \( s \) has pattern

\[
012 \ldots (k-2)(k-1)b, \text{ or } 012 \ldots (k-2)0b;
\]

and \( t \) has pattern

\[
012 \ldots (k-2)(k-1)b', \text{ or } 012 \ldots (k-2)0b'.
\]

And in any case, \( s \) and \( t \) differ in position \( k+1 \) and possibly in position \( k \).

The next proposition gives the pattern of \( \text{last}(\tilde{\gamma}_n | b) \) and \( \text{first}(\tilde{\gamma}_n | b) \) for an admissible suffix \( b \) in \( S_n \).

**Proposition 14.** Let \( b = b_{k+1} b_{k+2} \ldots b_n \) be an admissible suffix in \( S_n \). If \( s = \text{last}(\tilde{\gamma}_n | b) \), then the pattern of \( s \) is given by:

- if \( b_{k+1} = k \) or \( \sum_{i=k+1}^n b_i + (n-k) \) is odd, then
  \[
  s = 012 \ldots (k-2)(k-1)b;
  \]
- if \( b_{k+1} < k \) and \( \sum_{i=k+1}^n b_i + (n-k) \) is even, and either \( b_{k+1} = 0 \) or \( b_{k+1} \) is odd, then
  \[
  s = 012 \ldots (k-2)(\max\{0, b_{k+1} - 1\})b;
  \]
- if \( b_{k+1} < k \) and \( \sum_{i=k+1}^n b_i + (n-k) \) is even, and \( b_{k+1} > 0 \) is even, then
  \[
  s = 012 \ldots (k-3)(b_{k+1} - 2)(b_{k+1} - 1)b.
  \]

Similar results hold for \( s = \text{first}(\tilde{\gamma}_n | b) \) by replacing ‘odd’ by ‘even’, and vice versa, for the parity of \( \sum_{i=k+1}^n b_i + (n-k) \).

Proof. Let \( s = s_1 s_2 \ldots s_k b_{k+1} \ldots b_n = \text{last}(\tilde{\gamma}_n | b) \).

If \( b_{k+1} = k \) or \( \sum_{i=k+1}^n b_i + (n-k) \) is odd, then \( s_k \) is the largest value in \( \alpha_S(b) \), so \( s_k = k - 1 \), and by Remark 3, \( s_i = i - 1 \) for \( 1 \leq i \leq k \). So the first case holds.

If \( \sum_{i=k+1}^n b_i + (n-k) \) is even and \( b_{k+1} < k \), then \( s_k \) is the smallest value in \( \alpha_S(b) \), namely \( \max\{0, b_{k+1} - 1\} \), which is even if \( b_{k+1} = 0 \) or \( b_{k+1} \) is odd. Thus, by the definition of \( \prec \) order, \( s_{k-1} \) is the largest value in \( \alpha_S(s_k b) \), which is \( k - 2 \), and by Remark 3, the second case holds.

For the last case, as above, \( s_k = \max\{0, b_{k+1} - 1\} \), and considering \( b_{k+1} > 0 \) and even, it follows that \( s_k = b_{k+1} - 1 \) is odd. Thus \( s_{k-1} \) is the minimal value in \( \alpha_S(s_k b) \), that is \( b_{k+1} - 2 \), which in turns is even, and the last case holds.

The proof for the case \( s = \text{first}(\tilde{\gamma}_n | b) \) is similar.

**Theorem 4.** The list \( \tilde{\gamma}_n \) is 3-adjacent Gray codes.
Proof. Let $t = \text{succ}_{S}(s)$, and $k + 1$ be the rightmost position where $s$ and $t$ differ. Let us denote by $b$ the length $(n - k)$ suffix of $s$ and $b'$ that of $t$; so, $s = \text{last}(b | S_n)$ and $t = \text{first}(b' | S_n)$. It follows by Proposition 14 that $s_i = t_i = i - 1$ for all $i \leq k - 2$, and so the other differences possibly occur in position $k$ and in position $k - 1$.

Considering all valid combinations for $s$ and $t$ as given in Proposition 14, the proof of the adjacency is routine, and based on the following: $s_k \neq t_k$ if and only if $s_{k-1} \neq t_{k-1}$. It follows that $s$ and $t$ differ in one position, or three positions which are adjacent.

In addition, the lists $\tilde{A}_n$, $\tilde{R}_n$, and $\tilde{S}_n$, are circular Gray codes. This is a consequence of the following remarks based on Propositions 13 and 14:

- $\text{first}(\tilde{Y}_n) = 012 \ldots (n - 2)(n - 1)$;
- $\text{last}(\tilde{Y}_n) = 012 \ldots (n - 2)0$;

where $\tilde{Y}_n$ is one of the lists $\tilde{A}_n$, $\tilde{R}_n$, or $\tilde{S}_n$.

### 4.3 Generating algorithm for $\tilde{SE}_n$, $\tilde{A}_n$, $\tilde{R}_n$, and $\tilde{S}_n$

Here we explain algorithm $\text{Gen2}$ in Figure 3 which generates suffix partitioned Gray codes for restricted growth sequences; according to particular instances of the functions called by it, $\text{Gen2}$ produces the list $\tilde{SE}_n$, $\tilde{A}_n$, $\tilde{R}_n$, or $\tilde{S}_n$. Actually, for convenience, $\text{Gen2}$ produces length $(n + 1)$ sequences $s = s_1 s_2 \ldots s_{n+1}$ with $s_{n+1} = 0$, and so, neglecting the last value in each sequence $s$ the desired list is obtained. Notice that with this dummy value for $s_{n+1}$ we have $\mu_X(s_k s_{k+1} \ldots s_n) = \mu_X(s_k s_{k+1} \ldots s_n 0)$, for $k \leq n$, and similarly for $\alpha_X$.

In $\text{Gen2}$, the sequence $s$ is a global variable, and initialized by $01 \ldots (n - 1)0$, which is the first length $n$ sequence in $\prec_c$ order, followed by a $0$; and the main call is $\text{Gen2}(n + 1, 0, 0, 0)$. Procedure $\text{Gen2}$ has the following parameters (the first three of them are similar with those of procedure $\text{Gen1}$):

- $k$, the position in the sequence $s$ which is updated by the current call;
- $x$, the value to be assigned to $s_k$;
- $\text{dir}$, gives the direction in which $s_{k-1}$ covers $\alpha_X(s_k s_{k+1} \ldots s_n 0)$, the defining set of the current suffix;
- $v$, the value of $\mu_X(s_{k+1} \ldots s_n 0)$.

The functions called by $\text{Gen2}$ are given in Figures 4 and 5. They are principally based on the evaluation of $\alpha_X$ and $\mu_X$ for the current suffix of $s$, and are:

- $\text{Mu}_X(k, x, v)$ returns the value of $\mu_X(s_k s_{k+1} \ldots s_n 0)$, with $x = s_k$.
- $\text{IsDegreeOne}_X(k)$ stops the recursive calls when $\alpha_X(s_k s_{k+1} \ldots s_n 0)$ has only one element, namely $k - 2$. In this case, by Remark 3 the sequence is uniquely determined by the current suffix, and this prevents $\text{Gen2}$ to produce degree one calls. In addition, $\text{IsDegreeOne}_X(k)$ sets appropriately $d - 1$ values at the left of $s_k$, where $d$ is the upper bound of the Hamming distance in the list (changes at the left of $s_1$ are considered with no effect). This can be considered as a Path Elimination Technique [16].
• **Lowest\_X(k)** is called when **IsDegreeOne\_X** returns *false*, and gives the lowest value in \(\alpha_X(s_ks_{k+1}\ldots s_n0)\).

• **SecLargest\_X(k, u)**, is called when **IsDegreeOne\_X** returns *false*, and gives the second largest value in \(\alpha_X(s_ks_{k+1}\ldots s_n0)\) (the largest value being always \(k−2\)).

By this construction, algorithm **Gen2** has no degree one calls and it satisfies the CAT principle. Figure 6 shows the tree induced by the algorithm when generates \(\tilde{A}_4\).

```plaintext
procedure Gen2(k, x, dir, v)
  global n, s;
  s_k := x;
  u := Mu\_X(k, x, v);
  if IsDegreeOne\_X(k, u)
    then Print s;
  else
    c:=Lowest\_X(k);
    d:=SecLargest\_X(k, u);
    if dir mod 2 = 1
      then for i:=c to d do
        Gen2(k−1, i, i, u);
      enddo
    endif
    Gen2(k−1, k−2, (k−1−dir) mod 2, u);
    if dir mod 2 = 0
      then for i:=d downto c do
        Gen2(k−1, i, i + 1, u);
      enddo
    endif
  endif
end.
```

Figure 3: Algorithm **Gen2**, generating the list \(\tilde{X}_n\).
function Mu_A(m, i, w)
    if i ≥ w
        then return i;
    else if i ≥ s_{m+1}
        then return w;
            else return w - 1;
    endif
endif
end.

function IsDegreeOne_A(m, v)
    if v = m - 1 or
        (v = m - 2 and s_m = 0)
        then s_{m-1} := m - 2;
            return true;
        else return false;
    endif
end.

function Lowest_A(m)
    return 0;
end.

function SecLargest_A(m, w)
    if w = m - 2 and s_m > 0
        and s_m < m - 1
        then return s_m - 1;
            else return m - 3;
    endif
end.

function Mu_R(m, i, w)
    if i ≥ w
        then return i;
    else if s_{m+1} < w
        then return w;
            else return w - 1;
    endif
endif
end.

function IsDegreeOne_R(m, v)
    if v = m - 1 or
        (v = m - 2 and s_m < m - 2)
        then s_{m-1} := m - 2;
            return true;
        else return false;
    endif
end.

function Lowest_R(m)
    return 0;
end.

function SecLargest_R(m, w)
    return m - 3;
end.

(a) (b)

Figure 4: Particular functions called by Gen2, generating the lists: (a) \( \tilde{A}_n \), and (b) \( \tilde{R}_n \).
function Mu_S(m, i, w)
    return i;
end.

function IsDegreeOne_S(m, v)
    if s_m = m - 1
        then s_{m-1} := m - 2;
        s_{m-2} := m - 3;
        return true;
    else return false;
    endif
end.

function Lowest_S(m)
    if s_m > 1 and s_m ≤ m - 2
        then return s_m - 1;
        else return 0;
    endif
end.

function SecLargest_S(m, w)
    return m - 3;
end.

function Mu_SE(m, i, w)
    if m = n
        then return n - 1;
        else return w - 1;
    end.

function IsDegreeOne_SE(m, v)
    if m = 2
        then return true;
        else return false;
    end.

function Lowest_SE(m)
    return 0;
end.

function SecLargest_SE(m, w)
    return m - 3;
end.

(a) (b)

Figure 5: Particular functions called by Gen2, generating the lists: (a) \( \tilde{S}_n \), and (b) \( \tilde{SE}_n \).

Figure 6: The tree induced by the call Gen2(5, 0, 0, 0), generating the list \( \tilde{A}_4 \).
5 Final remarks

We conclude this paper by comparing for each of the sets $SE_n$, $A_n$, $R_n$ and $S_n$ the prefix partitioned Gray codes induced by $\prec$ order and the suffix partitioned one induced by $\prec_c$ order. This can be done by comparing the Hamming distance between all pairs of successive sequences either in the worst case or in average.

Table 1 summarizes Theorems 1, 3, and 4, and Proposition 2, and gives the upper bound of the Hamming distance (that is, the worst case Hamming distance) for the two order relations. It shows that for the sets $SE_n$ and $S_n$ these relations have same performances, and for the sets $A_n$ and $R_n$, $\prec_c$ order induces more restrictive Gray codes.

| Set | $\prec$ (RGC) order | $\prec_c$ (Co-RGC) order |
|-----|---------------------|-------------------------|
| $SE_n$ | 1                  | 1                        |
| $A_n$  | 3                  | 2                        |
| $R_n$  | 3                  | 2                        |
| $S_n$  | 3                  | 3                        |

Table 1: The bound of the Hamming distance between two successive sequences in $\prec$ and $\prec_c$ orders.

For a list $L$ of sequences, the average Hamming distance is defined as

$$\frac{\sum d(s, t)}{N - 1},$$

where the summation is taken over all $s$ in $L$, except its last element, $t$ is the successor of $s$ in $L$, $d$ is the Hamming distance, and $N$ the number of sequences in $L$.

Surprisingly, despite $\prec_c$ order has same or better performances in terms of worst case Hamming distance, if we consider the average Hamming distance, numerical evidences show that $\prec$ order is ‘more optimal’ than $\prec_c$ order on $A_n$, $R_n$ ($n \geq 5$), and $S_n$ ($n \geq 6$). And this phenomenon strengthens for large $n$; see Table 2.

| $n$ | $\prec$ (RGC) order | $\prec_c$ (Co-RGC) order |
|-----|---------------------|-------------------------|
| $SE_n$ | $A_n$ | $R_n$ | $S_n$ | $SE_n$ | $A_n$ | $R_n$ | $S_n$ |
| 4   | 1.21   | 1.21   | 1.31   | 1.14   | 1.14   | 1.15   |
| 5   | 1.13   | 1.12   | 1.29   | 1.19   | 1.18   | 1.24   |
| 6   | 1.09   | 1.07   | 1.27   | 1.23   | 1.20   | 1.31   |
| 7   | 1.06   | 1.06   | 1.26   | 1.25   | 1.22   | 1.35   |
| 8   | 1.04   | 1.04   | 1.25   | 1.26   | 1.23   | 1.37   |
| 9   | 1.03   | 1.03   | 1.24   | 1.28   | 1.24   | 1.39   |
| 10  | 1.02   | 1.03   | 1.23   | 1.28   | 1.24   | 1.41   |

Table 2: The average Hamming distance for $\prec$ order and $\prec_c$ order.

Algorithmically, $\prec_c$ order has the advantage that its corresponding generating algorithm, Gen2, is more appropriate to be parallelized than its $\prec$ order counterpart, Gen1. Indeed, the main call of Gen2 produces $n$ recursive calls (compare to two recursive calls produced by the
main call of \texttt{Gen1}), and so we can have more parallelized computations; and this is more suitable for large \( n \). See Figure 2 and 3 for examples of computational trees.

Finally, it will be of interest to explore order relation based Gray codes for restricted growth sequences defined by statistics other than those considered in this paper. In this vein we suggest the following conjecture, checked by computer for \( n \leq 10 \), and concerning descent sequences (defined similarly with ascent sequences in Section 2).

\textbf{Conjecture 1.} The set of length \( n \) descent sequences listed in \( \prec_c \) order is a 4-adjacent Gray code.
Appendix

| Sequence | $S_5$ | $R_5$ | $A_5$ |
|----------|-------|-------|-------|
| 00000    | ✓     | ✓     | ✓     |
| 00001    | ✓     | ✓     | ✓     |
| 00012    | ✓     | ✓     | ✓     |
| 00011    | ✓     | ✓     | ✓     |
| 00010    | ✓     | ✓     | ✓     |
| 00123    | ✓     | ✓     | ✓     |
| 00122    | ✓     | ✓     | ✓     |
| 00121    | ✓     | ✓     | ✓     |
| 00100    | ✓     | ✓     | ✓     |
| 01233    | ✓     | ✓     | ✓     |
| 01234    | ✓     | ✓     | ✓     |
| 01223    | ✓     | ✓     | ✓     |
| 01222    | ✓     | ✓     | ✓     |
| 01221    | ✓     | ✓     | ✓     |
| 01220    | ✓     | ✓     | ✓     |
| 01210    | ✓     | ✓     | ✓     |
| 01211    | ✓     | ✓     | ✓     |
| 01212    | ✓     | ✓     | ✓     |
| 01213    | ✓     | ✓     | ✓     |
| 01203    | ✓     | ✓     | ✓     |
| 01202    | ✓     | ✓     | ✓     |
| 01201    | ✓     | ✓     | ✓     |
| 01123    | ✓     | ✓     | ✓     |
| 01232    | ✓     | ✓     | ✓     |
| 01122    | ✓     | ✓     | ✓     |
| 01022    | ✓     | ✓     | ✓     |
| 01021    | ✓     | ✓     | ✓     |
| 01011    | ✓     | ✓     | ✓     |
| 01002    | ✓     | ✓     | ✓     |
| 01010    | ✓     | ✓     | ✓     |
| 01012    | ✓     | ✓     | ✓     |
| 00112    | ✓     | ✓     | ✓     |
| 00102    | ✓     | ✓     | ✓     |
| 00101    | ✓     | ✓     | ✓     |
| 00100    | ✓     | ✓     | ✓     |
| 01230    | ✓     | ✓     | ✓     |
| 01231    | ✓     | ✓     | ✓     |
| 00111    | ✓     | ✓     | ✓     |
| 01112    | ✓     | ✓     | ✓     |

Table 3: The sets $S_5$, $R_5$, and $A_5$ listed in $\prec$ order.

| Sequence | $S_5$ | $R_5$ | $A_5$ |
|----------|-------|-------|-------|
| 01234    | ✓     | ✓     | ✓     |
| 01233    | ✓     | ✓     | ✓     |
| 01023    | ✓     | ✓     | ✓     |
| 01123    | ✓     | ✓     | ✓     |
| 01223    | ✓     | ✓     | ✓     |
| 01213    | ✓     | ✓     | ✓     |
| 01013    | ✓     | ✓     | ✓     |
| 01203    | ✓     | ✓     | ✓     |
| 01202    | ✓     | ✓     | ✓     |
| 01102    | ✓     | ✓     | ✓     |
| 01002    | ✓     | ✓     | ✓     |
| 01012    | ✓     | ✓     | ✓     |
| 00122    | ✓     | ✓     | ✓     |
| 00112    | ✓     | ✓     | ✓     |
| 00112    | ✓     | ✓     | ✓     |
| 01232    | ✓     | ✓     | ✓     |

Table 4: The sets $S_5$, $R_5$, and $A_5$ listed in $\prec_c$ order.

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