Explicit MBR All-Symbol Locality Codes

Govinda M. Kamath, Natalia Silberstein, N. Prakash, Ankit S. Rawat, V. Lalitha, O. Ozan Koyluoglu, P. Vijay Kumar, and Sriram Vishwanath

Abstract

Node failures are inevitable in distributed storage systems (DSS). To enable efficient repair when faced with such failures, two main techniques are known: Regenerating codes, i.e., codes that minimize the total repair bandwidth; and codes with locality, which minimize the number of nodes participating in the repair process. This paper focuses on regenerating codes with locality, using pre-coding based on Gabidulin codes, and presents constructions that utilize minimum bandwidth regenerating (MBR) local codes. The constructions achieve maximum resilience (i.e., optimal minimum distance) and have maximum capacity (i.e., maximum rate). Finally, the same pre-coding mechanism can be combined with a subclass of fractional-repetition codes to enable maximum resilience and repair-by-transfer simultaneously.

I. BACKGROUND

A. Vector Codes

An $[n, K, d_{\text{min}}, \alpha]$ vector code over a field $\mathbb{F}_q$ is a code $C$ of block length $n$, having a symbol alphabet $\mathbb{F}_q^\alpha$ for some $\alpha > 1$, satisfying the additional property that given $c, c' \in C$ and $a, b \in \mathbb{F}_q$, $ac + be'$ also belongs to $C$. As a vector space over $\mathbb{F}_q$, $C$ has dimension $K$, termed the scalar dimension (equivalently, the file size) of the code and as a code over the alphabet $\mathbb{F}_q^\alpha$, the code has minimum distance $d_{\text{min}}$.

Associated with the vector code $C$ is an $\mathbb{F}_q$-linear scalar code $C^{(s)}$ of length $N = n\alpha$, where $C^{(s)}$ is obtained by expanding each vector symbol within a codeword into $\alpha$ scalar symbols (in some prescribed order). Given a generator matrix $G$ for the scalar code $C^{(s)}$, the first code symbol in the vector code is naturally associated with the $1$st code symbol $c_1$ as the $i$th thin column and to avoid confusion, the columns of $G$ themselves as thin columns.

B. Locality in Vector Codes

Let $C$ be an $[n, K, d_{\text{min}}, \alpha]$ vector code over a field $\mathbb{F}_q$, possessing a $(K \times n\alpha)$ generator matrix $G$. The $i$th code symbol, $c_i$, is said to have $(r, \delta)$ locality, $\delta \geq 2$, if there exists a punctured code $C_i := C|_{S_i}$ of $C$ (called a local code) with support $S_i \subseteq \{1, 2, \ldots, n\}$ such that

- $i \in S_i$,
- $|S_i| \leq n_L := r + \delta - 1$, and
- $d_{\text{min}}(C|_{S_i}) \geq \delta$.

The code $C$ is said to have $(r, \delta)$ information locality if there exists $l$ code symbols with $(r, \delta)$ locality and respective support sets $\{S_i\}_{i=1}^l$ satisfying

- $\text{Rank}(G|_{\bigcup_{i=1}^l S_i}) = K$.

The code $C$ is said to have $(r, \delta)$ all-symbol locality if all code symbols have $(r, \delta)$ locality. A code with $(r, \delta)$ information (respectively, all-symbol) locality is said to have full $(r, \delta)$ information (respectively, all-symbol) locality, if all local codes have parameters given by $|S_i| = r + \delta - 1$ and $d_{\text{min}}(C_i) = \delta$, for $i = 1, \ldots, l$.

The concept of locality for scalar codes, with $\delta = 2$, was introduced in [1] and extended in [2] and [3] to scalar codes with arbitrary $\delta$, and vector codes with $\delta = 2$, respectively. This was further extended to vector codes with arbitrary $\delta$ in [4] and [5], where, in addition to constructions of vector codes with locality, authors derive minimum distance upper bounds and also consider settings in which the local codes have regeneration properties.

Govinda M. Kamath, N. Prakash, V. Lalitha and P. Vijay Kumar are with the Department of ECE, Indian Institute of Science, Bangalore, India (email: {govinda, prakashnk, lalitha, vijay}@ece.iisc.ernet.in). Natalia Silberstein, Ankit S. Rawat, O. Ozan Koyluoglu, and Sriram Vishwanath are with Department of ECE, University of Texas at Austin, Austin, USA (email: {natalys, ankitsr, ozan, sriam}@austin.utexas.edu). This research is supported in part by the National Science Foundation under Grant 0964507 and in part by the NetApp Faculty Fellowship program. The work of V. Lalitha is supported by a TCS Research Scholarship.
Consider now a vector code \( C \) with full \((r, \delta)\) locality whose associated local codes \( C_i \) have parameters \([n_L, K_L, \delta]\). In this paper, we are interested in local codes that have the uniform rank accumulation property, in particular, local MBR codes and local fractional-repetition codes.

**Definition 1** (Uniform rank accumulation (URA) codes). Let \( G \) be a generator matrix for a code \( C \), and \( S_i \) be an arbitrary subset of \( i \) thick columns of \( G \), for some \( i = 1, \ldots, n \). Then, \( C \) is an URA code, if the restriction \( G|_{S_i} \) of \( G \) to \( S_i \), has rank \( \rho_i \) that is independent of the specific subset \( S_i \) of \( i \) indices chosen and given by \( \rho_i = \sum_{j=1}^{i} a_j \) for some set of non-negative integers \( \{a_j\} \).

We will refer to the sequence \( \{a_i, 1 \leq i \leq n\} \) as the rank accumulation profile of the code \( C \).

We now present the minimum distance upper bound given in \[4\] for the case when local codes are URA codes. Consider the finite length vector \((a_1, a_2, \ldots, a_{n_L})\), and its extension to a periodic semi-infinite sequence \( \{a_i\}_{i=1}^{\infty} \) of period \( n_L \) by defining \( a_{i+jn_L} = a_i \), \( 1 \leq i \leq n_L, \ j \geq 1 \). Let \( P(\cdot) \) denote the sequence of partial sums,

\[
P(s) = \sum_{i=1}^{s} a_i, \ s \geq 1.
\] (1)

Then, given integers \( u_1 \geq 0, \ 1 \leq u_0 \leq n_L, \ P(u_1n_L + u_0) = u_1K_L + P(u_0) \). Next, let us define the function \( P^{(\text{inv})} \) by setting \( P^{(\text{inv})}(\nu) \), for \( \nu \geq 1 \), to be the smallest integer \( s \) such that \( P(s) \geq \nu \). It can be verified that for \( v_1 \geq 0 \) and \( 1 \leq v_0 \leq K_L \),

\[
P^{(\text{inv})}(v_1K_L + v_0) = v_1n_L + P^{(\text{inv})}(v_0),
\]

where \( P^{(\text{inv})}(v_0) \leq r \) as \( 1 \leq v_0 \leq K_L \).

The minimum distance of a code \( C \) whose local codes \( C_i \) are URA codes can be bounded as follows.

**Theorem 1.1** (Theorem 5.1 of \[4\]). The minimum distance of \( C \) is upper bounded by

\[
d_{\min} \leq n - P^{(\text{inv})}(K) + 1.
\] (2)

The codes achieving the bound in (2) are referred to as codes having optimal locality. For such locality optimal codes, one can then analyze whether the code allows for efficient data storage in DSS. Towards this end, file size bound for codes with locality are given in \[5\] using the min-cut techniques similar to that of \[3\]. As noted in \[4\], when URA codes are used as local codes, the file size bound for \( d_{\min}\)-optimal codes can be represented in the form

\[
K \leq P(n - d_{\min} + 1) = \left( \left\lfloor \frac{n - d_{\min} + 1}{n_L} \right\rfloor - 1 \right) K_L + P(l_0),
\] (3)

where \( l_0 \in \{1, \ldots, n_L\} \) is such that

\[
n - d_{\min} + 1 = \left( \left\lfloor \frac{n - d_{\min} + 1}{n_L} \right\rfloor - 1 \right) n_L + l_0.
\]

We note that \( P(l_0) = P(r), \) for \( r \leq l_0 \).

C. MBR Codes

An \((n, k, d), (\alpha, \beta), K\) minimum-bandwidth regenerating (MBR) code is an \([n, K, d_{\min} = n - k + 1, \alpha]\) vector code satisfying additional constraints described below. The code is intended to be used in a distributed storage network in which each code symbol is stored within a distinct node. The code is structured in such a way that the entire file can be recovered by processing the contents of any \( k \), \( 1 \leq k \leq n \) nodes. Further, in case of a single node failure, the replacement node can reconstruct the data stored in the failed node by connecting to any \( d, k \leq d \leq n - 1 \), nodes and downloading \( \beta = \frac{a}{\alpha} \) symbols from each node. The scalar dimension (or file size) parameter \( K \) can be expressed in terms of the other parameters as:

\[
K = \left( dk - \binom{k}{2} \right) \beta.
\]
as proved in [6]. A cut-set bound derived from network coding shows us that the file size cannot be any larger, and thus, MBR codes are example of regenerating codes that are optimal with respect to file size. A regenerating code is said to be exact if the replacement of a failed node stores the same data as did the failed node, and functional otherwise. We are concerned here only with exact-repair codes. Constructions of MBR codes for all \( k \leq d = \alpha < n \) and \( \beta = 1 \) are presented in [7]. MBR codes with repair by transfer and \( d = n - 1 \) are presented in [8].

It can be inferred from the results in [8] that MBR codes are URA codes. In particular, for an \( ((n, k, d), (\alpha, \beta), K) \) MBR code, the rank accumulation profile is given by

\[
a_j = \begin{cases} 
\alpha - (j - 1)\beta, & 1 \leq j \leq k \\
0, & k + 1 \leq j \leq n.
\end{cases}
\] (4)

D. MBR-Local Codes

Let \( C \) be an \([n, K, d_{\text{min}}, \alpha] \) vector code with
- full \((r, \delta)\)-information locality with \( \delta \geq 2 \), and
- all of whose associated local codes \( C_i, i \in \mathcal{L} \) are MBR codes with identical parameters \((n_L = r + \delta - 1, r, d), (\alpha, \beta), K_L \)\).

Then, the dimension of each local code is given by

\[
K_L = \sum_{i=1}^{n_L} a_i = \alpha r - \binom{r}{2} \beta,
\] (5)

where \( \{a_i, 1 \leq i \leq n_L\} \) is the rank accumulation profile of the MBR code \( C \).

1) Minimum distance bound for MBR-Local Codes: As MBR codes are URA codes, from Theorem [11] we have

\[
d_{\text{min}} \leq n - P^{(\text{inv})}(K) + 1,
\] (6)

where, for MBR codes we have

\[
P^{(\text{inv})}(v_1 K_L + v_0) = v_1 n_L + \nu
\] (7)

for some \( v_1 \geq 0, 1 \leq v_0 \leq K_L \), and \( \nu \) is uniquely determined from \( \alpha(\nu - 1) - \binom{\nu - 1}{2} \beta < v_0 \leq \alpha \nu - \binom{\nu}{2} \beta \).

2) File size bound for MBR-Local Codes: From (3), the file size bound for an optimal locality code with MBR local codes is given by

\[
K \leq \left[ \left\lfloor \frac{n - d_{\text{min}} + 1}{n_L} \right\rfloor - 1 \right] K_L + \alpha \mu - \binom{\mu}{2} \beta,
\] (8)

where \( \mu = \min\{l_0, r\} \) with \( l_0 \) as defined in Subsection [1B]. Note that (8) follows from the rank accumulation profile of MBR codes, i.e., from (4).

E. Linearized Polynomials

A polynomial \( f(x) \) over the field \( \mathbb{F}_{q^m} \), is said to be linearized of \( q \)-degree \( t \), if

\[
f(x) = \sum_{i=0}^{t} u_i x^i, \quad u_i \in \mathbb{F}_{q^m}, \quad u_t \neq 0.
\] (9)

A linearized polynomial \( f(x) \) over \( \mathbb{F}_{q^m} \) satisfies the following property [9]:

\[
f(\lambda_1 \theta_1 + \lambda_2 \theta_2) = \lambda_1 f(\theta_1) + \lambda_2 f(\theta_2) \quad \forall \theta_1, \theta_2 \in \mathbb{F}_{q^m}, \lambda_1, \lambda_2 \in \mathbb{F}_q.
\] (10)

A linearized polynomial \( f(x) \) over \( \mathbb{F}_{q^m} \) of \( q \)-degree \( t, m > t \), is uniquely determined from its evaluation at a set of \((t + 1)\) points \( g_1, \ldots, g_{t+1} \in \mathbb{F}_{q^m} \), that are linearly independent over \( \mathbb{F}_q \).
F. Gabidulin Maximum Rank Distance Codes

Now, we present a construction of maximum rank distance codes, provided by Gabidulin in [10]. These codes can be viewed as a rank-metric analog of Reed-Solomon codes.

The rank of a vector \( \mathbf{v} \in \mathbb{F}_q^N \), denoted by \( \text{rank}(\mathbf{v}) \), is defined as the rank of the \( m \times N \) matrix \( \mathbf{V} \) over \( \mathbb{F}_q \), obtained by expansion of every entry of \( \mathbf{v} \) to a column vector in \( \mathbb{F}_q^m \), based on the isomorphism between \( \mathbb{F}_q^m \) and \( \mathbb{F}_q^N \). Similarly, for two vectors \( \mathbf{v}, \mathbf{u} \in \mathbb{F}_q^N \), the rank distance is defined by 
\[
d_R(\mathbf{v}, \mathbf{u}) = \text{rank}(\mathbf{V} - \mathbf{U}).
\]

An \([N, K, D]_{q^m}\) rank-metric code \( C \subseteq \mathbb{F}_q^N \) is a linear block code over \( \mathbb{F}_q \) of length \( N \), dimension \( K \) and minimum rank distance \( D \). A rank-metric code that attains the Singleton bound \( D \leq N - K + 1 \) in rank-metric is called a maximum rank distance (MRD) code. For \( m \geq N \), a construction of MRD codes, called Gabidulin codes, is given as follows [10].

A codeword in an \([N, K, D = N - K + 1]_{q^m}\) Gabidulin code \( C_{\text{Gab}}, m \geq N \), is defined as
\[
\mathbf{c} = (f(\theta_1), f(\theta_2), \ldots, f(\theta_N)) \in \mathbb{F}_{q^m}^N,
\]
where \( f(x) \) is a linearized polynomial over \( \mathbb{F}_{q^m} \) of degree at most \( K - 1 \) with the coefficients given by the information message, and where the \( \theta_1, \ldots, \theta_N \in \mathbb{F}_{q^m} \) are linearly independent over \( \mathbb{F}_q \) [10].

II. Construction of Codes with MBR Locality

In this section, we will present two constructions of codes with local regeneration. In both cases, the local codes are MBR codes with identical parameters and both codes are optimal, i.e., they achieve the upper bound of Theorem [11] on minimum distance. The first construction is an all-symbol locality construction, while the second has information locality.

The constructions presented in this paper, adopt the linearized polynomial approach made use of in [11], [12], [5]. In particular, similar to the constructions proposed in [12], [5], the constructions of this paper have a two-step encoding process with the first step utilizing Gabidulin codes, which in turn, are based on linearized polynomials. The first code construction given below also proves the tightness of the bound on minimum distance of codes with URA derived in [4] (Theorem 5.1) for the case when \( K_L \mid K \), where \( K_L \) is the scalar dimension of the local MBR code.

Consider a code \( C_{\text{Basic}} \) that is simply the concatenation of \( t \) local MBR codes having identical parameters \( ((n_L, k, d), (\alpha, \beta), K_L) \). Thus a typical codeword \( \mathbf{c} \in C_{\text{Basic}} \) looks like
\[
\mathbf{c} = (\mathbf{c}_1^{\text{mbr}}, \mathbf{c}_2^{\text{mbr}}, \ldots, \mathbf{c}_t^{\text{mbr}}),
\]
where each vector \( \mathbf{c}_i^{\text{mbr}} \) is a codeword belonging to the MBR code. The generator matrix \( G_{\text{Basic}} \) of the code will clearly have a block-diagonal structure. It is straightforward to show that the smallest number \( \rho \), such that any \( \rho \) thick columns of \( G_{\text{Basic}} \) have rank \( \geq K \) is given by \( P^{(\text{inv})}(K) \), for any \( 1 \leq K \leq tK_L \).

**Construction II.1.** We will describe the construction by showing how encoding of a message vector takes place. The encoding is illustrated in Fig. 4. Given the message vector \( \mathbf{u} \in \mathbb{F}_q^K \), we first encode \( \mathbf{u} \) to a \( tK_L \) long Gabidulin codeword using its \( tK_L \) linearly independent points (over \( \mathbb{F}_q \)) \{\( \theta_1, \theta_2, \ldots, \theta_{tK_L} \} \subset \mathbb{F}_q^m \), i.e., by applying an \([tK_L, K, tK_L - K + 1]_{q^m}\) Gabidulin code, assuming \( m \geq tK_L \). We then partition \( tK_L \) symbols of the Gabidulin codeword, \( (f(\theta_1), f(\theta_2), \ldots, f(\theta_{tK_L})) \), into \( t \) disjoint sets of \( K_L \) symbols each. Each of these sets is then fed in as a message vector to a bank of \( t \) identical MBR encoders whose outputs constitute \( ((n_L, r, d), (\alpha, \beta), K_L) \) MBR codes. If \( \{\mathbf{c}_i^{\text{mbr}} \mid i = 1, 2, \ldots, t\} \) is the resulting set of \( t \) codewords, these codewords are then concatenated to obtain the desired codeword \( \mathbf{c} \). The code \( C \) thus constructed has:

- length \( n = tn_L \)
- \( t \) local \( ((n_L, r, d), (\alpha, \beta), K_L) \) MBR codes with disjoint supports
- full \((r, \delta)\) all-symbol locality where \( \delta \) is defined from \( n_L = r + \delta - 1 \).

**Theorem II.2.** Given any set of parameters \( n, r, \delta, K \), such that \( n = tn_L \) and \( K \leq tK_L \), the construction II.1 yields an optimal MBR-local code with full \((r, \delta)\) all-symbol locality whose minimum distance is given by
\[
d_{\text{min}} = n - P^{(\text{inv})}(K) + 1.
\]
We first present a useful lemma on codes that are obtained by concatenating a Gabidulin code (over $\mathbb{F}_{q^m}$) with a vector code (over $\mathbb{F}_q$).

**Lemma II.3.** Let $G$ be the generator matrix of an $[n, J, d_{\text{min}}, \alpha]$ vector code over the field $\mathbb{F}_q$. Let $\tilde{J}$ be an integer such that $\tilde{J} \leq J$. Let $\rho$ be the smallest integer, such that the submatrix of $G$ obtained by selecting any $\rho$ thick columns of $G$ results in a matrix of rank $\geq \tilde{J}$. Let

$$\tilde{f}(x) = \sum_{i=0}^{\tilde{J} - 1} u_i x^i, u_i \in \mathbb{F}_{q^m}, \ m > J,$$

be a linearized polynomial of $q$-degree at most $\tilde{J} - 1$ over the extension field $\mathbb{F}_{q^m}$, for $m \geq J$. Let $\{\theta_i\}_{i=1}^J$ be any collection of $J$ elements of $\mathbb{F}_{q^m}$ that are linearly independent over $\mathbb{F}_q$. The mapping

$$(u_0, u_1, \cdots, u_{\tilde{J}-1}) \rightarrow (\tilde{f}(\theta_1), \tilde{f}(\theta_2), \cdots, \tilde{f}(\theta_J)) G$$

defines a linear code $C$ over $\mathbb{F}_{q^m}$ having message vector $(u_0, u_1, \cdots, u_{\tilde{J}-1})$. Then $C$ has minimum distance $D_{\text{min}}$ given by

$$D_{\text{min}} = n - \rho + 1,$$

i.e., $C$ is an $[n, \tilde{J}, D_{\text{min}}]$ code over $\mathbb{F}_{q^m}$.

**Proof:** Since $f(\cdot)$ is linearized, we can interchange linear operations with the operation of evaluation:

$$f((\theta_1, \theta_2, \cdots, \theta_j) G) = f(\tilde{f}(\theta_1, \theta_2, \cdots, \theta_j)) G.$$

We have extended here the definition of $f(\cdot)$ to vectors through termwise application. Consider next, the matrix product

$$\Gamma := [\theta_1 \theta_2 \cdots \theta_J] G.$$

In writing this, we have abused notation and identified elements in $\mathbb{F}_{q^m}$ with their representations as vectors over $\mathbb{F}_q$ lying in $\mathbb{F}_q^m$. The $m \times J$ matrix $[\theta_1 \theta_2 \cdots \theta_j]$ on the left has the property that all of its columns are linearly independent. Hence linear dependence relations amongst columns of $\Gamma$ are precisely those inherited from the matrix $G$. It follows that $\rho$ is also the smallest number, such that any $\rho$ thick columns of the product matrix $\Gamma$ have rank $\geq \tilde{J}$. Since $f(\cdot)$ is uniquely determined by its evaluation at a collection of $\tilde{J}$ linearly independent vectors lying in $\mathbb{F}_{q^m}$, it follows that the maximum number of erasures that the code $C$ can recover from is given by $n - \rho$. Then, we have

$$D_{\text{min}} = n - \rho + 1.$$
Remark 1. We note that whenever \( K = v_1 K_L + v_0, \) \( v_1 \geq 0, \) \( 1 \leq v_0 \leq K_L \) is such that \( v_0 = \nu \alpha - \binom{\nu}{2} \beta \) for some \( 1 \leq \nu \leq r, \) then the code constructed by Construction II.1 has maximum possible scalar dimension given in \( \text{[8]} \). This observation holds for the code we will construct using Construction II.4 as well.

Construction II.4. We describe here a method by which we construct a code of length \( n = t n_L + \Delta \). We then divide the first \( t K_L \) symbols of the Gabidulin codeword into \( t \) disjoint groups of equal size and encode each of these \( t \) groups using an \((n_L, r, d), (\alpha, \beta), K_L)\) MBR code (similar to the second step of encoding in Construction II.1). This gives us a code of length \( t n_L \) with MBR all-symbol locality, whose elements are \( \{c_i^{\text{mbr}} \mid i = 1, 2, \ldots, t\} \). We then partition the remaining \( \Delta \) symbols of the Gabidulin code into \( \Delta \) equal sets and denote the \( i^{th} \) set by \( c_{tn_L+i} \). The construction outputs \( (c_1^{\text{mbr}}, \ldots, c_t^{\text{mbr}}, c_{tn_L+1}, \ldots, c_{tn_L+\Delta}) \) as a final codeword. The resultant vector code \( C \) has:

- Length \( n = t n_L + \Delta \)
- \( t \) local \((n_L, r, d), (\alpha, \beta), K_L)\) MBR codes with disjoint support
- Full \((r, \delta)\) information locality

**Theorem II.5.** Given any set of parameters \( n, r, \delta, K \), such that \( n = t n_L + \Delta \) and \( K \leq t K_L \), Construction II.4 results in an optimal MBR-local code with \((r, \delta)\) information locality whose minimum distance is given by

\[
d_{\text{min}} = n - P^{(\text{inv})}(K) + 1.
\]

**Proof:** The proof follows along the same lines as the proof of Theorem II.2.

III. Fractional-Repetition Codes as Local Codes

In this section, we discuss the usage of fractional repetition (FR) codes as local codes in Constructions II.1 and II.4. FR codes can be viewed as a generalization of repair-by-transfer MBR codes, where a repair process is uncoded and table-based, i.e., FR codes have a “repair-by-transfer” property, while only specific sets of nodes of size \( d \) participate in a node repair process. For the sake of completeness, we provide an overview of the \( t \)-design-based construction for FR codes presented in [13].

Let \( t, n, w, \lambda \) be integers with \( n > w \geq t \) and \( \lambda > 0 \). A \( t-(n, w, \lambda) \) design is a collection \( \mathcal{B} \) of \( w \)-subsets (the blocks), of an \( n \)-set \( \mathcal{X} \) (the points), such that every \( t \)-subset of \( \mathcal{X} \) is contained in exactly \( \lambda \) blocks. Let \( x_1, \ldots, x_t \in \mathcal{X} \) be a set of \( t \) points. We denote by \( \lambda_s \) the number of blocks containing \( x_1, \ldots, x_s, 1 \leq s \leq t \). Then,

\[
\lambda_s = \binom{n-s}{t-s} \binom{t}{w-s} \binom{w-s}{t-s};
\]

the number of blocks in the \( t \)-design is \( b = \lambda_0 = \binom{n}{t} / \binom{w}{t} \); and each point in \( \mathcal{X} \) is contained in \( \lambda_1 \) blocks where \( \lambda_1 = \lambda(n-1) / (t-1) \) [9].

**Construction III.1.** Let \( B_1, \ldots, B_w \in \mathcal{B} \) be the blocks and \( x_1, \ldots, x_n \in \mathcal{X} \) the points of a \( t-(n, w, \lambda) \) design. Then the \( n \) nodes of a FR code \( C \) are given by the points of the design, i.e., a node \( N_i \) contains \( \alpha \triangleq \lambda_1 \) symbols given by \( N_i = \{j : x_i \in B_j\} \). Note that the cardinality of an intersection of any \( s \leq t \) nodes is given by the numbers \( \lambda_s \).

\[\text{The construction in [13] sets } t = 2 \text{ and } \lambda = 1; \text{ and the corresponding codes are called transposed codes.}\]
and hence the cardinality of a union of any $s \leq t$ nodes can be easily derived by the inclusion-exclusion formula. Let $k, K$ be two integers such that $k \leq t$ and

$$|\bigcup_{i=1}^{k-1} N_i| < K \leq |\bigcup_{i=1}^{k} N_i|. \quad \text{(12)}$$

Then we have an FR code over an alphabet of size $b$, with the property that there exists a set of $d$ nodes which can repair a failed node and from any set of $k$ nodes one can reconstruct the original $K$ symbols.

Given a message vector $[m_1 \ m_2 \ \cdots m_K]$, we encode the message symbols first by using an $[b, K, b-K+1]$ MDS code to produce $b$ coded symbols $(c_1, c_2 \cdots c_b)$ and then by employing the FR code based on the $t$-design to produce $n$ nodes each containing $\lambda_1$ symbols.

This family of FR codes based on $t$-designs is also an example of codes with uniform rank accumulation, and thus the bound of Theorem I.1 can be used here as well. Thus, we have the following result.

**Theorem III.2.** When FR codes based on a $t$-design obtained by Construction III.1 are used as the local codes in Constructions II.1 and II.4 then the resulting code with locality attains the bound of Theorem I.1 on minimum distance.

An example of an encoding is shown in Fig. 2 where the encoding is done using 2-$(7, 3, 1)$ design, also known as the Fano plane (see Fig. 3). When we replace a local MBR code with the FR code based on Fano plane in Fig. 1 we obtain a code with locality which has the optimal minimum distance.

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