Shortest unique palindromic substring queries in optimal time

Yuto Nakashima, Hiroe Inoue, Takuya Mieno, Shunsuke Inenaga, Hideo Bannai, Masayuki Takeda

1 Department of Informatics, Kyushu University
2 Japan Society for the Promotion of Science (JSPS), Japan

{yuto.nakashima,hiroe.inoue,takuya.mieno,inenaga,bannai,takeda}@inf.kyushu-u.ac.jp

Abstract

A palindrome is a string that reads the same forward and backward. A palindromic substring of a string \( S \) is called a shortest unique palindromic substring (SUPS) for an interval \([s, t]\) if \( S \) contains \([s, t]\), and every palindromic substring of \( S \) which contains interval \([s, t]\) is shorter than \( P \) occurs at least twice in \( S \). The SUPS problem is, given a string \( S \), to preprocess \( S \) so that for any subsequent query interval \([s, t]\) all the SUPSs for interval \([s, t]\) can be answered quickly. We present an optimal solution to this problem. Namely, we show how to preprocess a given string \( S \) of length \( n \) in \( O(n) \) time and space so that all SUPSs for any subsequent query interval can be answered in \( O(\alpha + 1) \) time, where \( \alpha \) is the number of outputs.

1 Introduction

A substring \( S[i..j] \) of a string \( S \) is called a shortest unique substring (SUS) for a position \( p \) if \( S[i..j] \) is the shortest substring t.s. \( S[i..j] \) is unique in \( S \) (i.e., \( S[i..j] \) occurs exactly once in \( S \)), and \( [i..j] \) contains \( p \) (i.e., \( i \leq p \leq j \)). Recently, Pei et al. \[13\] proposed the point SUS problem, preprocessing a given string \( S \) of length \( n \) so that we can return a SUS for any given query position efficiently. This problem was considered for some applications in bioinformatics, e.g., polymerase chain reaction (PCR) primer design in molecular biology. Pei et al. \[13\] proposed an algorithm which returns a SUS for any given position in constant time after \( O(n^2) \)-time preprocessing. After that, Tsuruta et al. \[15\] and Ileri et al. \[1] independently showed optimal \( O(n) \)-time preprocessing and constant query time algorithms. They also showed optimal \( O(n) \)-time preprocessing and \( O(k) \) query time algorithms which return all SUSs for any given position where \( k \) is the number of outputs. Moreover, Hon et al. \[6\] proposed an in-place algorithm which returns a SUS. A more general problem called interval SUS problem, where a query is an interval, was considered by Hu et al. \[7\]. They proposed an optimal \( O(n) \)-time preprocessing and \( O(k) \) query time algorithm which returns all SUSs containing a given query interval. Most recently, Mieno et al. \[12\] proposed an efficient algorithm for interval SUS problem when the input string is represented by run-length encoding.

In this paper, we consider a new variant of interval SUS problems concerning palindromes. A substring \( S[i..j] \) is called a palindromic substring of \( S \) if \( S[i..j] \) and the reversed string of \( S[i..j] \) is the same string. The study of combinatorial properties and structures on palindromes is still an important and well studied topic in stringology \[1, 3, 8, 14\]. Droubay et al. \[8\] showed a string of length \( n \) can contain at most \( n + 1 \) distinct palindromes. Moreover, Grault et al. \[5\] proposed a linear time algorithm for computing all distinct palindromes in a string.

Our new problem can be described as follows. A substring \( S[i..j] \) of a string \( S \) is called a shortest unique palindromic substring (SUPS) for an interval \([s, t]\) if \( S[i..j] \) is the shortest substring t.s. \( S[i..j] \) is unique in \( S \), \([i..j] \) contains \([s, t]\), and \( S[i..j] \) is a palindromic substring. The interval SUPS problem is to preprocess a given string \( S \) of length \( n \) so that we can return all SUPSs for any query interval efficiently. For this problem, we propose an optimal \( O(n) \)-time preprocessing and \( O(\alpha + 1) \)-time query algorithm, where \( \alpha \) is the number of outputs. Potential applications of our algorithm are in bioinformatics; It is known that the presence of particular (e.g., unique) palindromic sequences can affect immunostimulatory activities of oligonucleotides \[16\]. The size and the number of palindromes also influence the activity. Since any unique palindromic sequence can be obtained easily from a shorter unique palindromic sequences, we can focus on the shortest unique palindromic substrings.
The contents of our paper are as follows. In Section 2, we state some definitions and properties on strings. In Section 3, we explain properties on SUPS and our query algorithm. In Section 4, we show the main part of the preprocessing phase of our algorithm. Finally, we conclude.

2 Preliminaries

2.1 Strings

Let $\Sigma$ be an integer alphabet. An element of $\Sigma^*$ is called a string. The length of a string $S$ is denoted by $|S|$. The empty string $\varepsilon$ is a string of length 0, namely, $|\varepsilon| = 0$. Let $\Sigma^+$ be the set of non-empty strings, i.e., $\Sigma^+ = \Sigma^* - \{\varepsilon\}$. For a string $S = xyz$, $x$, $y$ and $z$ are called a prefix, substring, and suffix of $S$, respectively. A prefix $x$ and a suffix $z$ of $S$ are respectively called a proper prefix and proper suffix of $S$, if $x \neq S$ and $z \neq S$. The $i$-th character of a string $S$ is denoted by $S[i]$, where $1 \leq i \leq |S|$. For a string $S$ and two integers $1 \leq i \leq j \leq |S|$, let $S[i..j]$ denote the substring of $S$ that begins at position $i$ and ends at position $j$. For convenience, let $S[i..j] = \varepsilon$ when $i > j$.

2.2 Palindromes

Let $S^R$ denote the reversed string of $S$, that is, $S^R = S[|S|] \cdots S[1]$. A string $S$ is called a palindrome if $S = S^R$. Let $P \subset \Sigma^*$ be the set of palindromes. A substring $S[i..j]$ of $S$ is said to be a palindromic substring of $S$, if $S[i..j] \in P$. The center of a palindromic substring $S[i..j]$ of $S$ is $\lfloor \frac{i + j}{2} \rfloor$. Thus a string $S$ of length $n \geq 1$ has $2n - 1$ centers $(1, 1.5, \ldots, n - 0.5, n)$. The following lemma can be easily obtained by the definition of palindromes.

Lemma 1. Let $S$ be a palindrome. For any integers $i, j$ s.t. $1 \leq i \leq j \leq |S|$, $S[|S| - j + 1..|S| - i + 1] = S[i..j]^R$ holds.

2.3 MUPSs, SUPSs and our problem

For any non-empty strings $S$ and $w$, let $occ_S(w)$ denote the set of occurrences of $w$ in $S$, namely, $occ_S(w) = \{i \mid 1 \leq i \leq |S| - |w| + 1, w = S[i..i + |w| - 1]\}$. A substring $w$ of a string $S$ is called a unique substring (resp. a repeat) of $S$ if $|occ_S(w)| = 1$ (resp. $|occ_S(w)| \geq 2$). In the sequel, we will identify each unique substring $w$ of $S$ with its corresponding (unique) interval $[i, j]$ in $S$ such that $w = S[i..j]$. A substring $S[i..j]$ is said to be unique palindromic substring if $S[i..j]$ is a unique substring in $S$ and a palindromic substring. We will say that an interval $[i_1, j_1]$ contains an interval $[i_2, j_2]$ if $i_1 \leq i_2 \leq j_2 \leq j_1$ holds. The following notation is useful in our algorithm.

**Definition 2** (Minimal Unique Palindromic Substring (MUPS)). A string $S[i..j]$ is a MUPS in $S$ if $S[i..j]$ satisfies all the following conditions;

1. $S[i..j]$ is a unique palindromic substring in $S$,
2. $S[i + 1..j - 1]$ is a repeat in $S$ or $1 \leq |S[i..j]| \leq 2$.

Let $\mathcal{M}_S$ denote the set of intervals of all MUPSs in $S$ and let $mups_i = [b_i, e_i]$ denote the $i$-th MUPS in $\mathcal{M}_S$ where $1 \leq i \leq m$ and $m$ is the number of MUPSs in $S$. We assume that MUPSs in $\mathcal{M}_S$ are sorted in increasing order of beginning positions. For convenience, we define $mups_0 = [-1, -1], mups_{m+1} = [n + 1, n + 1]$.

**Example 3** (MUPS). For $S = acbaaabcbcbcaab$, $\mathcal{M}_S = \{[4, 6], [8, 12], [13, 16]\}$ (see also Fig. 7).

**Definition 4** (Shortest Unique Palindromic Substring (SUPS)). A string $S[i..j]$ is a SUPS for an interval $[s, t]$ in $S$ if $S[i..j]$ satisfies all the following conditions;

1. $S[i..j]$ is a unique palindromic substring in $S$,
2. $[i, j]$ contains $[s, t]$,
3. no unique palindromic substring $S[i'..j']$ containing $[s, t]$ with $j' - i' < j - i$ exists.

**Example 5** (SUPS). Let $S = acbaaabcbcbcaab$. SUPS for interval $[6, 7]$ is the $S[3..7] = bbaab$. SUPS for interval $[7, 8]$ are $S[2..8] = cbaaab$ and $S[7..13] = bcbcbcb$. SUPS for interval $[4, 13]$ does not exist. (see also Fig. 7).
In this paper, we tackle the following problem.

**Problem 6 (SUPS problem).**

- **Preprocess**: String $S$ of length $n$.
- **Query**: An interval $[s, t] (1 \leq s \leq t \leq n)$.
- **Return**: All the SUPSs for interval $[s, t]$.

### 2.4 Computation Model

Our model of computation is the word RAM: We shall assume that the computer word size is at least $\lceil \log_2 n \rceil$, and hence, standard operations on values representing lengths and positions of strings can be manipulated in constant time. Space complexities will be determined by the number of computer words (not bits).

### 3 Solution to the SUPS problem

In this section, we show how to compute all SUPSs for any query interval $[s, t]$.

#### 3.1 Properties on SUPS and MUPS

In our algorithm, we compute SUPSs by using MUPSs. Firstly, we show the following lemma. Lemma 7 states that MUPSs cannot nest in each other.

**Lemma 7.** For any pair of distinct MUPSs, one cannot contain the other.

**Proof.** Consider two MUPSs $u, v$ such that $u$ contains $v$. If $u$ and $v$ have the same center, then $u$ is not a MUPS. On the other hand, if $u$ and $v$ have a different center, we have from Lemma 1 and that $v$ is a palindromic substring, $v$ occurs in $u$ at least twice. This contradicts that $v$ is unique.

From this lemma, we can see that no pair of distinct MUPSs begin nor end at the same position. This fact implies that the number of MUPSs is at most $n$ for any string of length $n$. The following lemma states a characterization of SUPSs by MUPSs.

**Lemma 8.** For any SUPS $S[i..j]$ for some interval, there exists exactly one MUPS that is contained in $[i..j]$. Furthermore, the MUPS has the same center as $S[i..j]$.

**Proof.** Let $S[i..j]$ be a SUPS for some interval. $S[i..j]$ contains a MUPS $S[x_1..y_1]$ of the same center, i.e., $\frac{j-i}{2} = \frac{y_1-x_1}{2}$, s.t. $j - i \geq y_1 - x_1$. Suppose that there exists another MUPS $S[x_2..y_2]$ contained in $[i, j]$. From Lemma 7, $S[x_1..y_1]$ and $S[x_2..y_2]$ do not have the same center. On the other hand, if $S[x_1..y_1]$ and $S[x_2..y_2]$ have different centers, then $S[x_2..y_2]$ occurs at least two times in $S[i..j]$ by Lemma 1 since $S[x_2..y_2] = S[x_2..y_2]^R$. This contradicts that $S[x_2..y_2]$ is a MUPS.

---

Figure 1: This figure shows all MUPSs for $S = \text{acbaabcbcbaab}$ and some SUPS described in Example 5.
From the above lemma, any SUPS contains exactly one MUPS which has the same center (see also Fig. 1). Below, we will describe the relationship between a query interval \([s, t]\) and the MUPS contained in a SUPS for \([s, t]\). Before explaining this, we define the following notations.

- \(\mathcal{M}([s,t])\) : the set of MUPSs containing \([s,t]\).
- \(\text{predMUPS}[i] = i \text{ s.t. } i = \max\{k \mid e_k \leq t\}\).
- \(\text{succMUPS}[s] = i \text{ s.t. } i = \min\{k \mid s \leq b_k\}\).

In other words, \(\text{mups}_{\text{predMUPS}[i]}\) is the rightmost MUPS which ends before position \(t+1\), and \(\text{mups}_{\text{succMUPS}[s]}\) is the leftmost MUPS which begins after position \(s-1\).

**Lemma 9.** Let \(S[i..j]\) be a SUPS for an interval \([s, t]\). Then, the unique MUPS \(S[x..y]\) contained in \([i, j]\) is in \(\{\text{predMUPS}[i]\} \cup \mathcal{M}([s, t]) \cup \{\text{succMUPS}[s]\}\).

**Proof.** Assume to the contrary that there exists a SUPS \(S[i..j]\) that contains a MUPS \(S[x..y] \notin \{\text{predMUPS}[i]\} \cup \mathcal{M}([s, t]) \cup \{\text{succMUPS}[s]\}\). Since \(S[x..y] \notin \mathcal{M}([s, t])\), \([x, y]\) does not contain \([s, t]\). Thus, there can be the following two cases:

- If \(y < t\), there must exist MUPS \([x', y']\) s.t. \(y < y' \leq t\), since \(S[x, y] \neq \text{predMUPS}[i]\). By Lemma 7, \(x < x'\). Thus \(i \leq x < x' \leq y' \leq t \leq j\) holds. However, this contradicts Lemma 8.
- If \(s < x\), there must exist MUPS \([x', y']\) s.t. \(s \leq x' < x\), since \(S[x, y] \neq \text{succMUPS}[s]\). By Lemma 7, \(y' \leq y\). Thus \(i \leq s \leq x' \leq y' \leq y \leq j\) holds. However, this contradicts Lemma 8.

Therefore the lemma holds.

Next, we want to explain how SUPSs are related to MUPSs. It is easy to see that there may not be a SUPS for some query interval. We first show a case where there are no SUPSs for a given query. The following corollary is obtained from Lemma 8.

**Corollary 10.** Let \(S[x_1..y_1]\) and \(S[x_2..y_2]\) be MUPSs contained in a query interval \([s, t]\). There is no SUPS for an interval \([s, t]\).

From this corollary, a SUPS for an interval \([s, t]\) can exist if the number of MUPSs contained in \([s, t]\) is less than or equal to 1. The following two lemmas show what the SUPS for \([s, t]\) is, when \([s, t]\) contains only one MUPS, and when \([s, t]\) does not contain any MUPSs.

**Lemma 11.** Let \(S[x..y]\) be the only MUPS contained in the query interval \([s, t]\). If \(S[x-z, y+z]\) is a palindromic substring where \(z = \max(x-s, t-y)\), then \(S[x-z, y+z]\) is the SUPS for \([s, t]\). Otherwise, there is no SUPS for \([s, t]\).

**Proof.** Assume that there exists a SUPS \(u\) for \([s, t]\) which has the same center with a MUPS other than \(S[x..y]\). By the definition of SUPS, \(u\) should contain \([s, t]\). Since \([s, t]\) contains \([x, y]\), \(u\) contains two MUPSs, a contradiction. Thus, there can be no SUPS s.t. the center is not \(\frac{x+y}{2}\). It is clear that \(S[x-z, y+z]\) is a unique palindromic substring if \(S[x-z, y+z]\) is a palindromic substring where \(z = \max(x-s, t-y)\). Therefore the lemma holds.

**Lemma 12.** Let \([s, t]\) be the query interval. Then SUPSs for \([s, t]\) are the shortest of the following candidates.

1. \(S[x..y]\) s.t. \([x, y] \in \mathcal{M}([s, t])\),
2. \(S[x-t+y..t]\) s.t. \([x, y] = \text{predMUPS}([s, t])\), if it is a palindromic substring,
3. \(S[s..y+x-s]\) s.t. \([x, y] = \text{succMUPS}([s, t])\), if it is a palindromic substring.

**Proof.** It is clear that \(S[x..y]\) is a unique palindromic substring containing \([s, t]\) if \([x, y] \in \mathcal{M}([s, t])\) exists. It is also clear that if \([x, y] = \text{predMUPS}([s, t])\) or \([x, y] = \text{succMUPS}([s, t])\), then \(S[x-t+y..t]\) or \([s..y+x-s]\), respectively, are unique palindromic substrings, if they are palindromic substrings. By Lemma 9 we do not need to consider palindromic substrings which have the same center as MUPSs other than the candidates considered above. Thus the shortest of the candidates is SUPS for \([s, t]\) (see also Fig. 2).
From the above arguments, the number of MUPSs is useful to compute SUPSs for a query interval. The following lemma shows how to compute the number of MUPSs contained in a given interval.

**Lemma 13.** For any interval \([s, t]\),

- if \(\text{succMUPS}[s] > \text{predMUPS}[t]\), \([s, t]\) contains no MUPS,
- if \(\text{succMUPS}[s] = \text{predMUPS}[t]\), \([s, t]\) contains only one MUPS, \(\text{mups}_{\text{succMUPS}[s]} = \text{mups}_{\text{predMUPS}[t]}\), and
- if \(\text{succMUPS}[s] < \text{predMUPS}[t]\), \([s, t]\) contains at least two MUPSs.

**Proof.**

- Let \(j = \text{succMUPS}[s] > \text{predMUPS}[t] = i\). Then \(b_i < s \leq b_j\) and \(e_i \leq t < e_j\) hold, and thus neither of \(\text{mups}_i\) and \(\text{mups}_j\) are contained in \([s, t]\). If we assume that \([s, t]\) contains a MUPS \(\text{mups}_k\) for some \(k\), it should be that \(i < k < j\), \(b_i < s \leq b_k < b_j\). However, this contradicts that \(j = \text{succMUPS}[s]\) (see also the top in Fig. 3).

- Let \(\text{succMUPS}[s] = \text{predMUPS}[t] = i\). Since \(\text{succMUPS}[s] = i\), \(b_i - 1\) should be less than \(s\), and \(b_i\) at least \(s\). Since \(\text{predMUPS}[t] = i\), \(e_i + 1\) should be larger than \(t\), and \(e_i\) at most \(t\). Thus \([s, t]\) only contains \(\text{mups}_i\) (see also the middle in Fig. 3).

- Let \(i = \text{succMUPS}[s] < \text{predMUPS}[t] = j\). Then \(s \leq b_i < b_j\) and \(e_i < e_j \leq t\) hold, which implies \(s \leq b_i \leq e_i < t\) and \(s < b_j \leq e_j \leq t\). Thus, both \(\text{mups}_i\) and \(\text{mups}_j\) are contained in \([s, t]\) (see also the bottom in Fig. 3).

\(\Box\)

### 3.2 Tools

Here, we show some tools for our algorithm.

**Lemma 14** (e.g., [14]). For any interval \([i, j]\) in \(S\) of length \(n\), we can check whether \(S[i..j]\) is a palindromic substring or not in \(O(n)\) preprocessing time and constant query time with \(O(n)\) space.

Manacher’s algorithm [11] can compute all maximal palindromic substrings in linear time. If we have the array of radiuses of maximal palindromic substrings for all \(2n - 1\) centers, we can check whether a given substring \(S[i..j]\) is a palindromic substring or not in constant time.
3.2.1 Range minimum queries (RmQ)

Let $A$ be an integer array of size $n$. A range minimum query $RmQ_A(i, j)$ returns the index of a minimum element in the subarray $A[i, j]$ for given a query interval $[i, j][1 \leq i \leq j \leq n]$, i.e., it returns one of $\arg\min_{i \leq k \leq j} A[k]$. It is well-known (see e.g., [2]) that after an $O(n)$-time preprocessing over the input array $A$, $RmQ_A(i, j)$ can be answered in $O(1)$ time for any query interval $[i, j]$, using $O(n)$ space.

3.3 Algorithm

Due to the arguments in Section 3.1, if we can compute $predMUPS$, the shortest $MUPS$s in $\mathcal{M}(\{s, t\})$ and $succMUPS$ for a query interval $[s, t]$, then, we can compute $SUPS$s for $[s, t]$. Below, we will describe our solution to the $SUPS$ problem.
3.3.1 Preprocessing phase

First, we compute $\mathcal{M}_S$ for a given string $S$ of length $n$ in increasing order of beginning positions. We show, in the next section, that this can be done in $O(n)$ time and space. After computing $\mathcal{M}_S$, we compute the arrays $\text{predMUPS}$ and $\text{succMUPS}$. It is easy to see that we can also compute these arrays in $O(n)$ time by using $\mathcal{M}_S$. In the query phase, we are required to compute the shortest $\text{MUPS}$s that contain the query interval $[s, t]$. To do so efficiently, we prepare the following array. Let $\text{Mlen}$ be an array of length $m = |\mathcal{M}_S|$, and the $i$-th entry $\text{Mlen}[i]$ holds the length of $\text{mups}_i$, i.e., $\text{Mlen}[i] = |\text{mups}_i| = e_i - b_i + 1$. We also preprocess $\text{Mlen}$ for $\text{RmQ}$ queries. This can be done in $O(m)$ time and space as noted in Section 3.2. Thus, since $m = O(n)$, the total preprocessing is $O(n)$ time and space.

3.3.2 Query phase

First, we compute how many $\text{MUPS}$s are contained in a query interval $[s, t]$ by using Lemma 13, which we denote by $\text{num}$. This can be done in $O(1)$ time given arrays $\text{predMUPS}$ and $\text{succMUPS}$.

- If $\text{num} = 0$, let $\text{mups}_j = \text{predMUPS}([s, t])$ and $\text{mups}_j = \text{succMUPS}([s, t])$, i.e., $i = \text{predMUPS}[s]$ and $j = \text{succMUPS}[t]$. We check whether $S[b_i - t + e_i..t]$ and $S[s..e_j + b_j - s]$ are palindromic substrings or not. If so, then they are candidates of $\text{SUPS}$ for $[s, t]$ by Lemma 12. Let $q$ be the length of the shortest candidates which can be found in the above. Second, we compute the shortest $\text{MUPS}$ in $\mathcal{M}([s, t])$, if their lengths are at most $q$. In other words, we compute the smallest values in $\text{Mlen}[i + 1..j - 1]$, if they are at most $q$. We can compute all such $\text{MUPS}$s in linear time w.r.t. the number of such $\text{MUPS}$s by using $\text{RmQ}$ queries on $\text{Mlen}[i + 1..j - 1]$. Let $k = \text{RmQ}_{\text{Mlen}}(i + 1..j - 1)$ and $\text{Mlen}[k] \leq q$, then we consider the range $\text{Mlen}[i + 1..k - 1]$ and $\text{Mlen}[k + 1..j - 1]$ and recurse. Otherwise, we stop the recursion. Finally, we return the shortest candidates as $\text{SUPS}$.

- If $\text{num} = 1$, let $\text{mups}_j$ be the $\text{MUPS}$ contained in $[s, t]$. First, we check whether $S[b_i - z, e_i + z]$ is a palindromic substring or not by using Lemma 14 where $z = \max\{b_i - s, t - e_i\}$. If so, then return $[b_i - z, e_i + z]$, otherwise $\text{SUPS}$ for $[s, t]$ does not exist.

- If $\text{num} \geq 2$, then, from Corollary 10 $\text{SUPS}$ for $[s, t]$ does not exist.

Therefore, we obtain the following.

**Theorem 15.** After constructing an $O(n)$-space data structure of a given string of length $n$ in $O(n)$ time, we can compute all $\text{SUPS}$ for a given query interval $[s, t]$ in $O(\alpha + 1)$ time where $\alpha$ is the number of outputs.

4 Computing $\text{MUPS}$s

In this section, we show how to compute $\mathcal{M}_S$ in $O(n)$ time and space. Let $\text{DP}_S$ be the set of distinct palindromic substrings in $S$, and $\text{strM}_S = \{S[i..j] \mid [i..j] \in \mathcal{M}_S\}$. Our idea of computing $\mathcal{M}_S$ is based on the following lemma.

**Lemma 16.** $\text{strM}_S \subseteq \text{DP}_S$.

**Proof.** It is clear that any string in $\text{strM}_S$ is a palindromic substring of $S$. \qed

An algorithm for computing all distinct palindromic substrings in string in linear time and space was proposed by Groult et al. [5]. We show a linear time and space algorithm which computes $\mathcal{M}_S$ by modifying Groult et al.’s algorithm.

4.1 Tools

We show some tools for computing $\mathcal{M}_S$ below.

- **Longest previous factor array (LPF)** We denote the longest previous factor array of $S$ by $\text{LPF}_S$. The $i$-th entry ($1 \leq i \leq n$) is the length of the longest prefix of $S[i..n]$ which occurs at a position less than $i$.

- **Inverse suffix array (ISA)** We denote the inverse suffix array of $S$ by $\text{ISA}_S$. The $i$-th entry ($1 \leq i \leq n$) is the lexicographic order of $S[i..n]$ in all suffixes of $S$. 

7
• **Longest common prefix array (LCP)** We denote the longest common prefix array of \( S \) by \( LCP_S \). The \( i \)-th entry (\( 2 \leq i \leq n \)) is the length of the longest common prefix of the lexicographically \( i \)-th suffix of \( S \) and the \( (i-1) \)-th suffix of \( S \).

### 4.2 Computing distinct palindromes

Here, we show a summary of Groult et al.’s algorithm. The following lemma states the main idea.

**Lemma 17** \((\text{[3]})\). The number of distinct palindromic substrings in \( S \) is equal to the number of prefixes of \( S \) s.t. its longest palindromic suffix is unique in the prefix.

Since counting suffixes that uniquely occur in a prefix implies that only the leftmost occurrences of substrings, and thus distinct substrings are counted, their algorithm finds all the distinct palindromic substrings by:

- computing the longest palindromic suffix of each prefix of \( S \), and
- checking whether each longest palindromic suffix occurs uniquely in the prefix or not.

They first propose an algorithm which computes all the longest palindromic suffixes in linear time. They then check, in constant time, the uniqueness of the occurrence in the prefix by using the LPF array, thus computing \( DP_S \) in linear time and space.

### 4.3 Computing all MUPSs

Finally, we show how to modify Groult et al.’s algorithm. As mentioned, they compute the leftmost occurrence of each distinct palindromic substring. We call such a palindromic substring, the leftmost palindromic substring. It is clear that if a leftmost palindromic substring \( w \) is unique in \( S \) and is a minimal palindromic substring, then \( w \) is a **MUPS**. Thus, we add operations to check the uniqueness and minimality of each leftmost palindromic substring. We can do these operations by using **ISA** and **LCP** array.

Let \( S[i..j] \) be a leftmost palindromic substring in \( S \). First, we check whether \( S[i..j] \) is unique or not in \( S \). If \( ISA[i] = k \), \( S[i..n] \) is the lexicographically \( k \)-th suffix of \( S \). \( S[i..j] \) is unique in \( S \) if \( LCP[k] < j-i+1 \) and \( LCP[k+1] < j - i + 1 \). Thus we can check whether \( S[i..j] \) is unique or not in constant time. Finally, we check whether \( S[i..j] \) is a minimal palindromic substring or not. By definition, \( S[i..j] \) is minimal palindromic substring if \( j-i+1 \leq 2 \), i.e., \( S[i..j] \) has no shorter unique palindromic substring. If \( j-i+1 > 2 \), then we check whether \( S[i+1..j-1] \) is unique or not by using **ISA** and **LCP** in a similar way. Thus we can also check whether \( S[i..j] \) is minimal or not in constant time. By the above arguments, we can compute all **MUPSs** in linear time and space.

### 5 Conclusions

We consider a new problem called the shortest unique palindromic substring problem. We proposed an optimal linear time preprocessing algorithm so that all **SUPSs** for any given query interval can be answered in linear time w.r.t. the number of outputs. The key idea was to use palindromic properties in order to obtain a characterization of **SUPS**, more precisely, that a palindromic substring cannot contain a unique palindromic substring with a different center.

### References

[1] H. Bannai, T. Gagie, S. Inenaga, J. Kärkkäinen, D. Kempa, M. Piatkowski, S. J. Puglisi, and S. Sugimoto. Diverse palindromic factorization is NP-complete. In Proc. DLT 2015, pages 85–96, 2015.

[2] M. A. Bender and M. Farach-Colton. The LCA problem revisited. In Proceedings of the 4th Latin American Symposium on Theoretical Informatics, LATIN 2000, pages 88–94, 2000.

[3] X. Droubay, J. Justin, and G. Pirillo. Episturmian words and some constructions of de Luca and Rauzy. Theor. Comput. Sci., 255(1-2):539–553, 2001.
[4] G. Fici, T. Gagie, J. Kärkkäinen, and D. Kempa. A subquadratic algorithm for minimum palindromic factorization. *J. Discrete Algorithms*, 28:41–48, 2014.

[5] R. Groult, É. Prieur, and G. Richomme. Counting distinct palindromes in a word in linear time. *Inf. Process. Lett.*, 110(20):908–912, 2010.

[6] W.-K. Hon, S. V. Thankachan, and B. Xu. An in-place framework for exact and approximate shortest unique substring queries. In *ISAAC 2015*, pages 755–767, 2015.

[7] X. Hu, J. Pei, and Y. Tao. Shortest unique queries on strings. In *Proc. SPIRE 2014*, pages 161–172, 2014.

[8] T. I, S. Sugimoto, S. Inenaga, H. Bannai, and M. Takeda. Computing palindromic factorizations and palindromic covers on-line. In *Proc. CPM 2014*, pages 150–161, 2014.

[9] A. M. Ileri, M. O. Külekci, and B. Xu. Shortest unique substring query revisited. In *CPM 2014*, pages 172–181, 2014.

[10] E. Kuramoto, O. Yano, Y. Kimura, M. Baba, T. Makino, S. Yamamoto, T. Kataoka, and T. Tokunaga. Oligonucleotide sequences required for natural killer cell activation. *Japanese journal of cancer research*, 83(11):1128–1131, 1992.

[11] G. Manacher. A new linear-time “on-line” algorithm for finding the smallest initial palindrome of a string. *Journal of the ACM*, 22:346–351, 1975.

[12] T. Mieno, S. Inenaga, H. Bannai, and M. Takeda. Shortest unique substring queries on run-length encoded strings. In *Proc. MFCS 2016*, pages 69:1–69:11, 2016.

[13] J. Pei, W. C.-H. Wu, and M.-Y. Yeh. On shortest unique substring queries. In *Proc. ICDE 2013*, pages 937–948, 2013.

[14] M. Rubinchik and A. M. Shur. EERTREE: an efficient data structure for processing palindromes in strings. In *Combinatorial Algorithms - 26th International Workshop, IWOCA 2015, Verona, Italy, October 5-7, 2015, Revised Selected Papers*, pages 321–333, 2015.

[15] K. Tsuruta, S. Inenaga, H. Bannai, and M. Takeda. Shortest unique substrings queries in optimal time. In *Proc. SOFSEM 2014*, pages 503–513, 2014.

[16] S. Yamamoto, T. Yamamoto, T. Kataoka, E. Kuramoto, O. Yano, and T. Tokunaga. Unique palindromic sequences in synthetic oligonucleotides are required to induce IFN [correction of INF] and augment IFN-mediated [correction of INF] natural killer activity. *The Journal of Immunology*, 148(12):4072–4076, 1992.