Improved bounds on a generalization of Tuza’s conjecture

Abdul Basit Daniel McGinnis∗ Henry Simmons Matt Sinnwell Shira Zerbib†

Department of Mathematics
Iowa State University
Ames IA, U.S.A.

{abasit, dam1, has, mattds, zerbib}@iastate.edu

Submitted: Oct 26, 2021; Accepted: Sep 27, 2022; Published: Oct 21, 2022
© The authors. Released under the CC BY-ND license (International 4.0).

Abstract

For an $r$-uniform hypergraph $H$, let $\nu^{(m)}(H)$ denote the maximum size of a set $M$ of edges in $H$ such that every two edges in $M$ intersect in less than $m$ vertices, and let $\tau^{(m)}(H)$ denote the minimum size of a collection $C$ of $m$-sets of vertices such that every edge in $H$ contains an element of $C$. The fractional analogues of these parameters are denoted by $\nu^{*{(m)}}(H)$ and $\tau^{*{(m)}}(H)$, respectively. Generalizing a famous conjecture of Tuza on covering triangles in a graph, Aharoni and Zerbib conjectured that for every $r$-uniform hypergraph $H$, $\tau^{(r-1)}(H)/\nu^{(r-1)}(H) \leq \lceil \frac{r+1}{2} \rceil$. In this paper we prove bounds on the ratio between the parameters $\tau^{(m)}$ and $\nu^{(m)}$, and their fractional analogues. Our main result is that, for every $r$-uniform hypergraph $H$,

$$
\tau^{*{(r-1)}}(H)/\nu^{*{(r-1)}}(H) \leq \begin{cases} 
\frac{3r}{4} - \frac{r}{4(r+1)} & \text{for } r \text{ even}, \\
\frac{3r}{4} - \frac{r}{4(r+2)} & \text{for } r \text{ odd}.
\end{cases}
$$

This improves the known bound of $r-1$. We also prove that, for every $r$-uniform hypergraph $H$, $\tau^{(m)}(H)/\nu^{(m)}(H) \leq \text{ex}_r(n,m+1)$, where the Turán number $\text{ex}_r(n,k)$ is the maximum number of edges in an $r$-uniform hypergraph on $n$ vertices that does not contain a copy of the complete $r$-uniform hypergraph on $k$ vertices. Finally, we prove further bounds in the special cases $(r,m) = (4,2)$ and $(r,m) = (4,3)$.

Mathematics Subject Classifications: 05C65, 05D15

∗Supported by NSF grant DMS-1839918 (RTG).
†Supported by NSF grant DMS-1953929.
1 Introduction

1.1 Definitions

We restrict our attention to r-uniform hypergraphs (or r-graphs), i.e., hypergraphs where all edges are of size r. Denote by $H[X]$ the sub-hypergraph of $H$ induced by $X \subseteq V(H)$, that is, containing all edges that are contained in $X$. Throughout the paper, we set $V(H) = [n]$, and use the abbreviation $a_1a_2\ldots a_k$ to denote the set $\{a_1, a_2, \ldots, a_k\} \subseteq V(H)$. We shall also identify $H$ with its edge set $E(H)$.

A matching in a hypergraph $H$ is a set of disjoint edges. The matching number $\nu(H)$ is the maximum size of a matching in $H$. A cover is a set of vertices intersecting all edges of $H$. The covering number $\tau(H)$ is the minimum size of a cover. Clearly, in an r-graph $H$ we have $\nu(H) \leq \tau(H) \leq r\nu(H)$.

More generally, an m-matching $M$ in a hypergraph $H$ is a set of edges such that any two edges in $M$ intersect in less than $m$ vertices. An m-cover $C$ of $H$ is a collection of m-sets of vertices such that every edge in $H$ contains at least one member of $C$. So a 1-matching is a matching and a 1-cover is a cover. Denote by $\nu^{(m)}(H)$ the m-matching number, that is, the maximum size of an m-matching in $H$ and by $\tau^{(m)}(H)$ the m-covering number, namely the minimum size of an m-cover of $H$.

For a hypergraph $H$, let $H^{(m)}$ be the hypergraph whose vertex set is $\binom{V(H)}{m}$ and whose edge set is $\{\binom{e}{m} : e \in H\}$. Note that $M \subset H$ is an m-matching in $H$ if and only if $\{\binom{e}{m} : e \in M\}$ is a matching in $H^{(m)}$. Additionally, a collection $C$ of m-sets is an m-cover of $H$ if and only if $C$ is a cover of $H^{(m)}$. Thus, $\nu^{(m)}(H) = \nu(H^{(m)})$ and $\tau^{(m)}(H) = \tau(H^{(m)})$.

We will also consider fractional variants of these parameters. A fractional matching in $H$ is a function $s : H \rightarrow \mathbb{R}_{\geq 0}$ satisfying $\sum_{e \ni v} s(e) \leq 1$ for every $v \in V(H)$. The size of a fractional matching is $|s| = \sum_{e \ni H} s(e)$, and the fractional matching number of $H$, denoted by $\nu^*(H)$, is the maximum size of a fractional matching of $H$. Similarly, a fractional cover of $H$ is a function $t : V(H) \rightarrow \mathbb{R}_{\geq 0}$ satisfying $\sum_{e \ni v} t(v) \geq 1$ for every $e \in H$. The size of a fractional cover is $|t| = \sum_{v \in V(H)} t(v)$, and the fractional covering number of $H$, denoted by $\tau^*(H)$, is the minimum size of a fractional cover of $H$.

Determining the fractional matching and covering numbers is a linear programming problem. The two problems form a dual pair so by the LP duality principle, $\tau^*(H) = \nu^*(H)$ for every hypergraph $H$. Thus for an r-graph $H$ we have

$$\nu(H) \leq \nu^*(H) = \tau^*(H) \leq \tau(H) \leq r\nu(H). \quad (1)$$

Define a fractional m-matching and fractional m-cover in $H$ to be a fractional matching and fractional cover in $H^{(m)}$, respectively, and let $\nu^{*(m)}(H) = \nu^*(H^{(m)})$ and $\tau^{*(m)}(H) = \tau^*(H^{(m)})$. Then (1) implies

$$\nu^{(m)}(H) \leq \nu^{*(m)}(H) = \tau^{*(m)}(H) \leq \tau^{(m)}(H) \leq \left(\frac{r}{m}\right) \nu^{(m)}(H).$$
1.2 Tuza’s conjecture and its generalizations

For a finite graph $G$, let $\nu_t(G)$ be the maximum size of a set of edge-disjoint triangles in $G$, and let $\tau_t(G)$ is the minimum size of a set $C$ of edges with the property that each triangle contains a member of $C$. Clearly, $\nu_t(G) \leq \tau_t(G) \leq 3\nu_t(G)$. A famous conjecture of Tuza states:

**Conjecture 1** (Tuza [21]). For any graph $G$, $\tau_t(G) \leq 2\nu_t(G)$.

The inequality is seen to be sharp by taking $G$ to be $K_4$ or $K_5$ (or a disjoint union of these), and is close to sharp in other cases (see e.g. [2, 8]). The best known general bound is $\tau_t(G) \leq \frac{66}{23} \nu_t(G)$, due to Haxell [7].

Note that if $T(G)$ is the hypergraph whose edges are triples of vertices forming a triangle in $G$, then $\tau_t(G) = \tau^{(2)}(T(G))$ and $\nu_t(G) = \nu^{(2)}(T(G))$, and thus Tuza’s conjecture states that $\nu^{(2)}(T(G)) \leq 2\tau^{(2)}(T(G))$ for every graph $G$. In this paper we are interested in a generalization of Conjecture 1 proposed by Aharoni and Zerib [1]. They conjectured the same inequality holds for any 3-graph.

**Conjecture 2** (Aharoni-Zerib [1, Conjecture 1.2]). For any 3-graph $H$,

$$\tau^{(2)}(H) \leq 2\nu^{(2)}(H).$$

They also conjectured that a similar phenomenon should hold much more generally.

**Conjecture 3** (Aharoni-Zerib [1, Conjecture 1.10]). Fix $r \geq 3$. For any $r$-graph $H$,

$$\tau^{(r-1)}(H) \leq \left[ \frac{r + 1}{2} \right] \nu^{(r-1)}(H).$$

Conjecture 3 is based on a more general conjecture involving functions $h$ and $g$, defined ahead, which states that $h(r, m) = g(r, m)$ for all $r, m$ along with the observation that, for an $r$-graph $H$, if $\nu^{(r-1)}(H) = 1$ then $\tau^{(r-1)} \leq \left[ \frac{r+1}{2} \right]$. Let $\mathcal{H}_r$ be the family of all $r$-graphs, for some $r \geq 3$. For an integer $2 \leq m \leq r$, we will be interested in the following functions (first defined in [1]):

- $h(r, m) = \sup \left\{ \frac{\tau^{(m)}(H)}{\nu^{(m)}(H)} : H \in \mathcal{H}_r \right\}$,
- $g(r, m) = \sup \left\{ \tau^{(m)}(H) : H \in \mathcal{H}_r \text{ and } \nu^{(m)}(H) = 1 \right\}$,
- $h^*(r, m) = \sup \left\{ \frac{\tau^{(m)}(H)}{\nu^{(m)}(H)} : H \in \mathcal{H}_r \right\}$,
- $g^*(r, m) = \sup \left\{ \tau^{(m)}(H) : H \in \mathcal{H}_r \text{ and } \nu^{(m)}(H) = 1 \right\}$,
- $j^*(r, m) = \sup \left\{ \frac{\tau^{(m)}(H)}{\nu^{(m)}(H)} : H \in \mathcal{H}_r \right\}$.
The following proposition is easily proved from the definitions:

**Proposition 4.** For any integers $2 \leq m \leq r$, the following holds.

(a) $g^*(r, m) \leq g(r, m) \leq h(r, m),

(b) $h^*(r, m) \leq h(r, m),

(c) $j^*(r, m) \leq h(r, m).

Given a graph $G$, let $H(G, r)$ be the $r$-graph whose vertex set is $V(G)$ and whose edges are cliques of size $r$ in $G$. Let $\mathcal{G}_r \subseteq \mathcal{H}_r$ be the family of all $r$-graphs obtained in this manner. For each of the functions defined above, we add $\circ$ in the subscript to denote the function where $\mathcal{H}_r$ is replaced by $\mathcal{G}_r$. For example

$$g_\circ(r, m) = \sup \{ \tau^{(m)}(H) : H \in \mathcal{G}_r \text{ and } \nu^{(m)}(H) = 1 \}.$$ 

Observe that $f_\circ(r, m) \leq f(r, m)$, where $f$ is any of the functions $h, g, h^*, g^*, j^*$.

Using the notation above, Conjectures 1, 2 and 3 can be stated as:

**Conjecture 5.**

(a) $h_\circ(3, 2) \leq 2$ (*Tuza*),

(b) $h(3, 2) \leq 2$ (*Aharoni-Zerbib*),

(c) $h(r, r - 1) \leq \lceil \frac{r+1}{2} \rceil$ for $r \geq 3$ (*Aharoni-Zerbib*).

In this paper we determine, or give bounds on the values of the functions $h, g, h^*, g^*, j^*, h_\circ, g_\circ, h^*_\circ, g^*_\circ, j^*_\circ$, for certain sets of parameters $r, m$.

### 1.3 Our results and organization of the paper.

In Section 2 we focus on the case $m = r - 1$. That is, we consider the functions $g(r, r - 1)$, $h(r, r - 1)$, and their variants defined above. In [10] Krivelevich proved that $h^*_\circ(3, 2) \leq 2$. Aharoni and Zerbib [1] generalized this to $r$-graphs, and proved for any $r \geq 3$,

$$h^*(r, r - 1) \leq r - 1. \quad (2)$$

This follows from the observation that if $H$ is an $r$-graph then $H^{(r-1)}$ cannot contain the $r$-uniform projective plane $\mathbb{P}_r$, and a theorem of Füredi [5] stating that if an $r$-graph $H$ does not contain $\mathbb{P}_r$, then $\tau^*(H) \leq (r - 1)\nu(H)$. For an introduction to projective planes, see e.g. [17].

Here we improve the bound (2) for all $r \geq 4$. Moreover, in the case $r = 3$ we give a new (and shorter) constructive proof of (2), which does not use Füredi’s theorem.

**Theorem 6.** $h^*(3, 2) = 2$, and $h^*(4, 3) \leq \frac{8}{3}$. 

---

**The Electronic Journal of Combinatorics** 29(4) (2022), #P4.14
Theorem 7. For every \( r \geq 5 \),

\[
h^\ast(r, r-1) \leq \begin{cases} \frac{3}{4} r - \frac{r}{4(r+1)} & \text{for } r \text{ even,} \\ \frac{3}{4} r - \frac{r}{4(r+2)} & \text{for } r \text{ odd.} \end{cases}
\]

For \( r \geq 6 \), it is also possible to further improve Theorem 7, but the optimization is much more involved, and leads to a negligible gain for large \( r \). For example, when \( r = 6 \), it is possible to obtain an upper bound of \( 17/4 \) (vs. \( 30/7 \) implied by Theorem 7).

In Section 3, we study the case \((r, m) = (4, 2)\). Some bounds for these parameters and certain families of 4-graphs \( H \in \mathcal{G}_4 \) are given in Szttapolow [18, Chapter 5], e.g., when the corresponding graph \( G \) is the complete graph, for 4-partite graphs, and planar graphs. In [1] it was shown that \( g(4, 2) = 4 \). Completing the picture, we prove

Theorem 8.

(a) \( g^\ast(4, 2) = 2.5 \),

(b) \( g_o(4, 2) = 3 \),

(c) \( g^*(4, 2) = 3.5 \).

Observe, in particular, that \( g(4, 2) \neq g_o(4, 2) \), unlike the case \( r = 3 \) where it is known that \( g(3, 2) = g_o(3, 2) \) (see [1]) and conjectured that \( h(3, 2) = h_o(3, 2) \).

In [19] it was proved that \( h^o(4, 2) \leq 4.5 \) (the more general bound \( h^*(4, 2) \leq 4.5 \) was obtained in [1]). We improve this bound:

Theorem 9. \( h^o(4, 2) \leq 4 \).

Finally, in Section 4, we turn our attention to the function \( j^\ast \). In [1], it was shown that \( j^*(3, 2) = 2 \) and \( j^*(4, 2) \leq 4 \). Guruswami and Sandeep [6] made significant progress, proving that \( j^*(r, 2) \leq r^2/4 \) and \( j^*(r, r-1) \leq r/2 + \sqrt{2r \ln r} \). For the more general case, they give the bound \( j^*(r, m) = c(m)(r^m) \) where \( c(m) \to 1/2 + o(1) \) as \( m \to r - 1 \). In their paper, they also observed a connection to the hypergraph Turán problem. Here we exploit this connection more explicitly and bound \( j^*(r, m) \) in terms of hypergraph Turán numbers.

Let \( ex_r(n, k) \) be the maximum number of edges in an \( r \)-graph on \( n \) vertices that does not contain a copy of \( K^r_k \), the complete \( r \)-graph on \( k \) vertices. Turán [20] determined the value of \( ex_2(n, k) \) and posed the problem of determining the limit \[
\pi(k, r) = \lim_{n \to \infty} \frac{ex_r(n, k)}{\binom{n}{r}},
\]
for \( 2 < r < k \). This has proven to be a notoriously hard problem and even the first non-trivial case \( r = 3 \) and \( k = 4 \) remains open. For a survey of the problem and related results see [9, 14].

Theorem 10. \( j^*(r, m) \leq ex_m(r, m + 1) \).
When \( m = r - 1 \), we obtain the bound of \( j^*(r, r - 1) \leq r - 1 \) (this bound also appears in [6]). For \( m = 2 \), Theorem 10, together with Mantel’s Theorem [12], recovers the bounds \( j^*(3, 2) \leq 2 \), \( j^*(4, 2) \leq 4 \) and \( j^*(r, 2) \leq \frac{r^2}{4} \). More generally, Theorem 10 together with known results about hypergraph Turán numbers can be used to obtain explicit values. For example, paired with results of Chung and Lu [3], Markström [13] and Sidorenko [15], respectively, we obtain

\[
\begin{align*}
\bullet \quad & \lim_{r \to \infty} \frac{j^*(r, 3)}{\binom{r}{3}} \leq \frac{3 + \sqrt{17}}{12} \approx 0.5936. \\
\bullet \quad & \lim_{r \to \infty} \frac{j^*(r, 4)}{\binom{r}{4}} \leq \frac{1753}{2380} \approx 0.73655. \\
\bullet \quad & \lim_{r \to \infty} \frac{j^*(r, m)}{\binom{r}{m}} \leq 1 - \frac{1}{m}.
\end{align*}
\]

2 The case \( m = r - 1 \)

2.1 Preliminaries

Let \( H \in \mathcal{H}_r \) be an \( r \)-graph, and \( M \) be a maximum \((r - 1)\)-matching in \( H \). We say an edge \( e \) in \( H \) is of type-\( i \) (with respect to \( M \)), for \( 1 \leq i \leq r \), if \( e \) intersects exactly \( i \) edges of \( M \) in \( r - 1 \) vertices each, and \( e \) intersects every other edge of \( M \) in at most \( r - 2 \) vertices. Let \( T_i \subseteq H \) denote the set of edges of type \( i \). For \( e \in M \), let \( H(e) = \{f \in H : |f \cap e| \geq r - 1\} \), \( T_i(e) = T_i \cap H(e) \) and \( H_i(e) = T_i(e) \cup \{e\} \). For \( x \in \binom{r}{i-1} \), if there exists an edge \( f \in T_i(e) \) such that \( f \cap e = x \), then we say that \( x \) is an indispensable \((r - 1)\)-set. When \( r = 4 \), then we refer to an indispensable 3-set as an indispensable triple.

Observation 11. Let \( H \in \mathcal{H}_r \), and let \( M \) be a maximum \((r - 1)\)-matching in \( H \) with \( e \in M \), and let \( f, g \in T_i(e) \). Then \( |f \cap g| \geq r - 1 \). In other words, \( \nu^{(r-1)}(H_1(e)) = 1 \).

Proof. Assume to the contrary that \( |f \cap g| < r - 1 \). Then \((M \setminus \{e\}) \cup \{f, g\}\) is an \((r - 1)\)-matching of size greater than \(|M|\), a contradiction. \( \square \)

The following lemma gives a structural characterization of \( H_1(e) \).

Lemma 12. For \( H \in \mathcal{H}_r \), let \( M \) be a maximum \((r - 1)\)-matching in \( H \). Let \( e \in M \) be an edge such that \( T_1(e) \neq \emptyset \). Then one of the following holds:

(a) \( \bigcap H_1(e) = r - 1 \). If this is true, we set \( p(e) = \bigcap H_1(e) \).

(b) There exists a vertex \( v \in V(H) \setminus e \) such that \( v \in f \) for each \( f \in T_1(e) \). If such a vertex exists, we denote it by \( v(e) \).

Proof. We may assume \( |T_1(e)| > 1 \), otherwise the claim is trivial. For an edge \( f \in T_1(e) \) let \( f = x \cup \{v_f\} \), where \( v_f \notin e \). Let \( f, g \) be distinct edges in \( T_1(e) \). By Observation 11, if \( Y = f \cap g \) then \( |Y| = r - 1 \).
Assume first $|Y \cap e| = r - 1$. Let $h \in T_1(e)$, and assume $h$ does not contain $Y$. Then $|h \cap f| = r - 1$ and $|h \cap g| = r - 1$ implies $\{v_f, v_g\} \subset h$, implying $|h \cap e| \leq r - 2$, a contradiction.

Otherwise we have $|Y \cap e| < r - 1$. Then $|f \cap g| = r - 1$ implies $v_f = v_g$. Similarly, every edge $h \in T_1(e)$ has both $|f \cap h \cap e| < r - 1$ and $|g \cap h \cap e| < r - 1$, and thus $v_h = v_f$. □

Remark 13. If $|T_1(e)| > 1$, then exactly one of (a) or (b) holds. If (b) holds then $|T_1(e)|$ equals the number of indispensable $(r - 1)$-sets (namely, every indispensable $(r - 1)$-set belongs to exactly one edge in $T_1(e)$), and moreover $|e \cap f \cap g| = r - 2$ for distinct edges $f, g \in T_1(e)$.

Before proving Theorem 7, we first give the following weaker bound as a warm-up. When $r = 2$, the following proposition recovers the optimal bound $h^*(2, 1) = 1.5$ (see e.g. [11]).

Proposition 14. For any $r \geq 2$,

$$\frac{r + 1}{2} \leq h^*(r, r - 1) \leq \frac{3}{4} r.$$

Proof. To see the lower bound, consider the hypergraph $H = \binom{[r+1]}{r}$. Clearly $\nu^{(r-1)}(H) = 1$. The function $t : \binom{[r+1]}{r-1} \to \mathbb{R}$ with $t(x) = 1/r$ is a fractional $(r - 1)$-cover of size $(r + 1)/2$, and the function $s : H \to \mathbb{R}$ with $s(e) = 1/2$ is a fractional $(r - 1)$-matching of size $(r + 1)/2$. This implies $\tau^*(H) = \nu^*(H) = (r + 1)/2$.

For the upper bound, let $H \in \mathcal{H}_r$ be an $r$-graph with $\nu^{(r-1)}(H) = k$, and let $M$ be a maximum $(r - 1)$-matching in $H$. Define the function $t : \binom{[r+1]}{r-1} \to \mathbb{R}$ by $t(x) = 1/2$ if $x \in \{\binom{e}{r-1} : e \in M\}$, and $t(x) = 0$ otherwise. Note that $t$ is a fractional $(r - 1)$-cover of the type-$i$ edges for $i > 1$ and the edges in $M$. It remains to cover type-1 edges, each of which already has weight 1/2 by $t$.

Let $e \in M$. We define a function $t_e : \binom{[r+1]}{r-1} \to \mathbb{R}$ as follows. If $e$ satisfies Lemma 12(a), we set $t_e(p(e)) = 1/2$. If $e$ satisfies Lemma 12(b), then we set $t(x) = \frac{1}{2(r-1)}$ for $(r - 1)$-sets $x$ satisfying $|x \cap e| = r - 2$ and $v(e) \in x$. For any other $(r - 1)$-set we set $t_e(x) = 0$.

It is now easy to check that the function $t + \sum_{e \in M} t_e$ is a fractional $(r - 1)$-cover of size at most $\frac{3}{4} r |M|$. □

2.2 Proofs of Theorems 6 and 7

Let $H \in \mathcal{H}_r$ be an $r$-graph, and let $M$ be a maximum $(r - 1)$-matching in $H$. For $e \in M$, recall that an indispensable $(r - 1)$-set $x \in \binom{e}{r-1}$ is an $(r - 1)$-set such that there exists a type-1 edge $f \in H$ with $f \cap e = x$. We refer to $f$ as a witness of indispensability of $x$ in $e$. For $0 \leq i \leq r$, let $M_i$ be the set of all edges in $M$ containing exactly $i$ indispensable $(r - 1)$-sets. For $e, f \in M$ and a type-2 edge $h$, we say that $h$ connects $e$ and $f$ if $|e \cap h| = |f \cap h| = r - 1$ (namely, $e, f$ are the two edges in $M$ witnessing the fact that $h$ is type-2). We define $M^+ = M_3$ if $r = 3$, and $M^+ = M_{r-1} \cup M_r$ if $r \geq 4$. Let $M^- = M \setminus M^+$. 

THE ELECTRONIC JOURNAL OF COMBINATORICS 29(4) (2022), #P4.14 7
Lemma 15. Let $e \in M^+$. Let $f \in M$ and suppose there exists an edge connecting $e$ and $f$. Then there exists $g \in T_1(e)$ such that:

(a) $|g \cap h| < r - 1$ for any $h$ connecting $e$ and $f$,

(b) $|g \cap a| < r - 1$ for each $a \in T_1(f)$.

Proof. Let $h$ be an edge connecting $e$ and $f$. Note first that $|e \cap f| = r - 2$ and $e \cap f \subset h \subset e \cup f$. Since $e$ has at least 3 indispensable $(r - 1)$-sets at most two of which contain $e \cap f$, there exists an indispensable $(r - 1)$-set $x$ in $e$ such that $e \cap f \not\subseteq x$. Let $g \in T_1(e)$ be a witness for $x$. Then $g$ satisfies the required properties. Properties (a) and (b) follow from the fact, proved ahead, that the vertex $v := g \setminus e$ is not contained in $f$ (here we slightly abuse notation and identify the singleton set with the vertex). Indeed $v \notin f$ together with $e \cap f \not\subseteq g$ and, for any $h'$ connecting $e$ and $f$, $h' \subseteq e \cup f$ implies that $|g \cap h'| \leq r - 2$, so (a) is satisfied. We also have $|g \cap f| = r - 3$ implying $|g \cap a| \leq r - 2$ for every $a \in T_1(f)$, so (b) is satisfied.

To see that $v \notin f$, let $y$ be an indispensable $(r - 1)$-set in $e$ such that $e \cap f \subset y$ and let $a \in T_1(e)$ be a witness for $y$. Such an $(r - 1)$-set must exist since $e$ has at least $r - 1$ indispensable $(r - 1)$-sets at most $r - 2$ of which do not contain $e \cap f$. Lemma 12 applies and by Remark 13, $a$ contains the vertex $v$. Now, if $v \in f$ then $|a \cap f| \geq r - 1$, a contradiction. □

Almost immediately, we obtain the following.

Corollary 16. Let $e, f \in M^+$. Then there is no type-2 edge connecting $e$ and $f$.

Proof. Suppose that $g$ connects $e$ and $f$. By Lemma 15, there exist edges $a \in T_1(e)$ and $b \in T_1(f)$ such that $|a \cap g| < r - 1$, $|b \cap g| < r - 1$, and $|a \cap b| < r - 1$. It follows that $(M \setminus \{e, f\}) \cup \{a, b, g\}$ is an $(r - 1)$-matching of size greater than $|M|$, a contradiction. □

We call a type-2 edge bad if it connects an edge $e \in M^+$ and an edge $f \in M^-$. If a type-2 edge is not bad, then it is good. For an edge $e \in M_i$, let $B(e)$ be the set of all bad type-2 edges connecting $e$ to other edges in $M$.

Corollary 17. Let $e \in M^- \setminus M_0$, and suppose $g \in B(e)$ connects $e$ and $f \in M^+$. Then $|g \cap h| \geq r - 1$ for every $h \in T_1(e)$.

Proof. By Lemma 15, there exists an edge $a \in T_1(f)$ such that $|a \cap g| < r - 1$ and $|a \cap h| < r - 1$ for every $h \in T_1(e)$. If there exists $h \in T_1(e)$ such that $|h \cap g| < r - 1$, then $(M \setminus \{e, f\}) \cup \{a, g, h\}$ is an $(r - 1)$-matching of size greater than $|M|$, a contradiction. □

Lemma 18. Let $e \in M^- \setminus (M_0 \cup M_1)$, and suppose $g \in B(e)$ connects $e$ and $f \in M^+$. Then the following hold:

(a) $v(e) = g \setminus e$

(b) If $x$ in an indispensable $(r - 1)$-set in $e$, then $e \cap f \not\subseteq x$. 

THE ELECTRONIC JOURNAL OF COMBINATORICS 29(4) (2022), #P4.14
Proof. Let $x$ be an indispensable $(r - 1)$-set in $e$ such that $|x \cap g| = r - 2$, which exists because $e \in M_i$ with $i \geq 2$, and let $a \in T_1(e)$ be an edge witnessing $x$. By Corollary 17, $|a \cap g| \geq r - 1$ implying that $a$ contains the vertex $(g \setminus e) \in f$. By Lemma 12, there exists a vertex $v(e) \notin e$ such that $v$ is contained in every edge of $T_1(e)$. It follows that $v := v(e) = (g \setminus e) \in f$ (as before, we do not distinguish between the singleton set and the vertex).

Now suppose that an edge $b \in T_1(e)$ contains $e \cap f$. Since $v \in b$, we have $|b \cap f| \geq r - 1$, a contradiction. It follows that no indispensable $(r - 1)$-set in $e$ contains $e \cap f$. \hfill $\square$

**Corollary 19.** If $r = 3$, then there are no bad edges connecting $e \in M_2$ and $f \in M_3$.

**Proof.** Suppose $g$ connects $e \in M_2$ and $f \in M_3$. This implies $|e \cap f| = 1$. Since $e$ has two indispensable triples, at least one of them contains $e \cap f$, contradicting Lemma 18. \hfill $\square$

**Corollary 20.** Suppose $r \geq 4$, and let $e \in M^- \setminus (M_0 \cup M_1)$ such that $B(e) \neq \emptyset$. Let $i$ be such that $e \in M_i$. Then there exist at most $\binom{r}{2} - i(r - 1) + \binom{i}{2}$ $(r - 1)$-sets such that every edge in $B(e)$ contains at least one of them.

**Proof.** By Lemma 18(a), any edge $g$ that connects $e$ and $f \in M^+$ must contain $e \cap f$ and $v(e)$. Since $e \cap f$ is not contained in any indispensable $(r - 1)$-sets (by Lemma 18(b)), it suffices to bound the number of $(r-2)$-sets of $e$ that are not contained in any indispensable $(r - 1)$-sets of $e$. By inclusion-exclusion, this number is at most $\binom{r}{2} - i(r - 1) + \binom{i}{2}$. \hfill $\square$

**Remark 21.** If $r \geq 4$, then, for any $e \in M_{r-2}$, there exists an $(r - 1)$-set contained in all edges of $B(e)$.

**Lemma 22.** If $e \in M_1$, then one of the following holds:

(a) all edges in $T_1(e)$ and $B(e)$ share one $(r - 1)$-set $w(e)$, or

(b) $|T_1(e)| = 1$.

**Proof.** Let $x$ be the indispensable $(r-1)$-set in $e$, and suppose there exists an edge $g \in B(e)$ that connects $e$ and $f \in M^+$ such that $x \notin g$. Let $h \in T_1(e)$ be a witness for $x$. By Corollary 17, we have $|h \cap g| \geq r - 1$ implying that $h$ contains the vertex $g \setminus e$, i.e., $h$ is determined uniquely. It follows that $|T'(e)| = 1$. \hfill $\square$

We are now ready to prove Theorems 6 and 7. Since the proofs are similar, we only give details for the more involved proof of Theorem 7. For Theorem 6, we simply describe the fractional cover and leave the verification as an exercise.

**Proof of Theorem 7.** Let $\alpha := \alpha(r) = \frac{r+2}{2(r+1)}$ for $r$ even, and $\alpha := \alpha(r) = \frac{r+3}{2(r+2)}$ for $r$ odd. For $r \geq 5$, let $H \in \mathcal{H}_e$ and let $M$ be a maximum $(r - 1)$-matching in $H$. For every $e \in M$, we define $t_e : \binom{V(H)}{r-1} \to \mathbb{R}$ as follows (if we do not explicitly specify $t_e(x)$ for some $x \in \binom{V(H)}{r-1}$, then $t_e(x) = 0$):

1. If $e \in M_0$: Set $t_e(x) = \alpha$ for every $x \in \binom{V(H)}{r-1}$.
2. If $e \in M_1$: We define functions $t^0_e, t^1_e : (V(H))_{r-1} \to \mathbb{R}$ and set $t_e(x) = t^0_e(x) + t^1_e(x)$.

Set $t^0_e(x) = 1/2$ for every $x \in (r_{-1})$. If $e$ satisfies Lemma 22(a), then set $t^1_e(w(e)) = 1/2$; otherwise set $t^1_e(x) = \frac{1}{2r}$ for each $x \in (r_{-1})$ where $f$ is the unique element of $T_1(e)$.

3. If $e \in M_i$ for $2 \leq i \leq r - 3$: Set $t_e(x) = \alpha$ for every $x \in (r_{-1})$. Lemma 12(b) applies and by Remark 13 there are exactly $i$ edges in $T_1(e)$. Observation 11 implies that there exist at most $\lceil i/2 \rceil$ $(r-1)$-sets such that every edge in $T_1(e)$ contains one of these $(r-1)$-sets. For each such $(r-1)$-set $y$, set $t_e(y) = 1 - \alpha$.

4. If $e \in M_{r-2}$: We define functions $t^0_e, t^1_e, t^2_e : (V(H))_{r-1} \to \mathbb{R}$ and set $t_e(x) = t^0_e(x) + t^1_e(x) + t^2_e(x)$.

Set $t^0_e(x) = 1/2$ for every $x \in (r_{-1})$. Lemma 12(b) applies and by Remark 13 there are exactly $r - 2$ edges in $T_1(e)$. By Observation 11 there exist $\lceil \frac{r-2}{2} \rceil$ $(r-1)$-sets such that every edge in $T_1(e)$ contains one of them. For each such $(r-1)$-set $y$, set $t^1_e(y) = 1/2$.

By Lemma 18, if $B(e) \neq \emptyset$ then all edges in $B(e)$ contain an $(r-1)$-set $z$. Set $t^2(z) = \alpha - 1/2$.

5. If $e \in M^+$: Set $t_e(x) = 1 - \alpha$ for every $x \in (r_{-1})$. For each $(r-1)$-set $y$ that contains $v(e)$ and $r-2$ vertices from $e$, set $t_e(y) = \frac{\alpha}{r-1}$.

Let $t = \sum_{e \in M} t_e$. It is easy to check that $t$ covers all edges in $M$, type-1 edges, and type-$i$ edges for $3 \leq i \leq r$. Note that $\alpha \geq 1/2$ for all $r$, implying that good type-2 edges are covered. If a type-2 edge is in $B(e)$ for $e \in M_0$, then it is also easily seen to be covered. A type-2 edge in $B(e)$ for $e \in M_1$ gets weight at least $\frac{1}{2} + (1 - \alpha) = \frac{2r+1}{2^r+2}$ from $(r-1)$-sets contained in edges of $M$, and an additional weight $\frac{1}{2r}$ from $t^1_e$ (which is sufficient by Lemma 22). If a type-2 edge is in $B(e)$ for $e \in \bigcup_{i=3}^r M_i$, then it receives weight at least 1 from $(r-1)$-sets contained in edges of $M$. If a type-2 edge is in $B(e)$ for $e \in M_{r-2}$, then it receives weight at least $1/2 + 1 - \alpha$ from $(r-1)$-sets contained in edges of $M$, and $\alpha - 1/2$ from $t^2_e$.

Finally, we have

$$|t_e| \leq \begin{cases} r\alpha & \text{if } e \in M_0 \\ \frac{r+1}{2} & \text{if } e \in M_1 \\ r\alpha + (1 - \alpha) \left\lceil \frac{r-3}{2} \right\rceil & \text{if } e \in \bigcup_{i=3}^r M_i \\ \frac{r}{2} + \frac{1}{2} \left\lceil \frac{r-2}{2} \right\rceil + \alpha - \frac{1}{2} & \text{if } e \in M_{r-2} \\ r(1 - \alpha) + \left(\frac{r}{2}\right)^{\alpha-1} & \text{if } e \in M^+ \end{cases}$$

In each case, it is easy to check that the asserted bound holds.

Proof of Theorem 6. Let $H \in H_3$, and let $M$ a maximum 2-matching in $H$. For every $e \in M$, we define $t_e : (V(H))_2 \to \mathbb{R}$ as follows (if we do not explicitly specify $t_e(x)$ for some $x \in (V(H))_2$, then $t_e(x) = 0$):
1. If $e \in M_0$: Set $t_e(x) = 2/3$ for every $x \in \binom{e}{2}$.

2. If $e \in M_1$: We define functions $t^0_e, t^1_e : \binom{V(H)}{2} \to \mathbb{R}$ and set $t_e(x) = t^0_e(x) + t^1_e(x)$.
   
   Set $t^0_e(x) = 1/2$ for every $x \in \binom{e}{2}$.
   
   If $e$ satisfies Lemma 22(a), then set $t^1_e(w(e)) = 1/2$; otherwise set $t^1_e(x) = 1/6$ for each $x \in \binom{e}{2}$ where $f$ is the unique element of $T_1(e)$.

3. If $e \in M_2$: Set $t_e(x) = 1/2$ for every $x \in \binom{e}{3}$. By Observation 11, the two edges in $T_1(e)$ must intersect in a pair $y$. Set $t_e(y) = 1/2$.

4. If $e \in M_3$: Set $t_e(x) = 1/3$ for every $x \in \binom{e}{3}$.

Then $t = \sum_{e \in M} t_e$ is a cover of size at most $2|M|$.

Now let $H \in \mathcal{H}_4$, and let $M$ be a maximum 3-matching in $H$. For every $e \in M$, we define $t_e : \binom{V(H)}{3} \to \mathbb{R}$ as follows (if we do not explicitly specify $t_e(x)$ for some $x \in \binom{V(H)}{3}$, then $t_e(x) = 0$):

1. If $e \in M_0$: Set $t_e(x) = 2/3$ for every $x \in \binom{e}{3}$.

2. If $e \in M_1$: We define functions $t^0_e, t^1_e : \binom{V(H)}{3} \to \mathbb{R}$ and set $t_e(x) = t^0_e(x) + t^1_e(x)$.
   
   Set $t^0_e(x) = 1/2$ for every $x \in \binom{e}{3}$.
   
   If $e$ satisfies Lemma 22(a), then set $t^1_e(w(e)) = 1/2$; otherwise set $t^1_e(x) = 1/6$ for each $x \in \binom{e}{3}$ where $f$ is the unique element of $T_1(e)$.

3. If $e \in M_2$: We define functions $t^0_e, t^1_e : \binom{V(H)}{3} \to \mathbb{R}$ and set $t_e(x) = t^0_e(x) + t^1_e(x)$.
   
   Set $t^0_e(x) = 1/2$ for every $x \in \binom{e}{3}$. Lemma 12 applies and by Remark 13 there are exactly two edges in $T_1(e)$ which, by Observation 11, must intersect in a triple $y$. Set $t^0_e(y) = 1/2$.
   
   By Corollary 20, if $B(e) \neq \emptyset$ then all edges in $B(e)$ contain a triple $z$. Set $t^1_e(z) = 1/6$.

4. If $e \in M^+$: Set $t_e(x) = 1/3$ for every $x \in \binom{e}{3}$. Since there are at most four edges in $T_1(e)$ any two of which intersect in a triple, there are two triples $x$ and $y$ such that any element of $T_1(e)$ contains either $x$ or $y$. Set $t_e(x) = t_e(y) = 2/3$.

Then $t = \sum_{e \in M} t_e$ is a cover of size at most $\frac{8}{3}|M|$.  \qed

### 3 Fractional 2-covers in 4-graphs

#### 3.1 Preliminaries

Let $H \in \mathcal{H}_4$, and let $M$ be a maximum 2-matching in $H$. If we also have $H \in \mathcal{G}_4$ then edges of $H$ correspond to $K_4$’s in the graph $G$ with $V(G) = V(H)$. To avoid confusion, throughout this section by edges we will always mean edges of $H$ and will refer only to subgraphs of $G$. 
We say an edge in $H$ is of type-1 (with respect to $M$) if it intersects one edge of $M$ in at least two vertices and shares at most one vertex with every other edge of $M$. Let $T_1 \subseteq H$ denote the set of type-1 edges. For $e \in M$, let $H(e) = \{ f \in H : |f \cap e| \geq 2 \}$ and $T_1(e) = T_1 \cap H(e)$. Note that type-1 edges behave similarly to type-1 edges in Section 2 giving the following analogue of Observation 11.

**Observation 23.** Let $H \in \mathcal{H}_1$, let $M$ be a maximum 2-matching in $H$ with $e \in M$, and let $f, g \in T_1(e)$. Then $|f \cap g| \geq 2$.

For $e \in M$, a pair $p \in \binom{e}{2}$ is indispensable in $e$ if there exists a type-1 edge $f$ with $e \cap f = p$ and we refer to $f$ as a witness of indispensability of $p$. The following is an immediate consequence of Observation 23.

**Corollary 24.** Let $e \in M$. If there exist two disjoint indispensable pairs $p_1, p_2$ in $e$ with witnesses $f_1, f_2$ respectively, then $f_1 \cap f_2$ is a pair $q = q(p_1, p_2)$ disjoint from $e$.

If we also have $H \in \mathcal{G}_4$, then $G[e \cup q]$ is the graph $K_6$ implying that $H[e \cup q] = \binom{e\cup q}{4}$.

**Example 25.** Let $H \in \mathcal{G}_4$ be the complete 4-graph on six vertices, i.e., $H$ is the set of all $K_4$’s contained in $K_6$. Then $\nu(2)(H) = 1$, and $\tau(2)(H) = \nu(2)(H) = 2.5$.

Note that Example 25 is maximal in the sense that, for a 4-graph $H'$ with at least seven edges containing a copy of $H$, we have $\nu(2)(H') > 1$. Hence, Corollary 24 plays a key role in understanding the structure of $H \in \mathcal{G}_4$.

### 3.2 Proof of Theorem 8 (a) and (b)

The lower bound follows from Example 25.

For the upper bound, let $H \in \mathcal{G}_4$ with $\nu(2)(H) = 1$ and let $e \in H$. Since $\nu(2)(H) = 1$, every edge of $H$ must share at least two vertices with $e$. Note that if $p_1, p_2 \in \binom{e}{2}$ are such that $p_1 \cap p_2 = \emptyset$, then $\{p_1, p_2\}$ covers all edges that share at least three vertices with $e$. It follows that if there is at most one indispensable pair, then $\tau(2)(H) \leq 2$. Indeed, we may take an indispensable pair along with the (unique) pair disjoint from it to be the cover.

By Corollary 24, if $p_1, p_2$ are disjoint indispensable pairs in $e$ then $H[e \cup q(p_1, p_2)] = \binom{e\cup q}{4}$, implying that $H = \binom{e\cup q}{4}$, since otherwise $\nu(2)(H) > 1$. It follows that $\tau(2)(H) \leq 3$ and $\tau(2)(H) \leq 2.5$.

From here on we assume that there are at least two indispensable pairs, no two of which are disjoint. In particular, there are at most three indispensable pairs. Without loss of generality, let $e = 1234$, and $p_1 = 12, p_2 = 13$ be indispensable pairs with witnesses $f_1 = 1256$ and $f_2$ respectively. Since $f_2$ must intersect $f_1$ in at least two vertices, it contains at least one of the vertices 5 or 6. Without loss of generality, assume that $f_2 = 135u$ where $u \in V(H) \setminus 24$. It suffices to consider the following cases.

**Case 1:** $p_1$ and $p_2$ are the only indispensable pairs. Notice that any edge that intersects $e$ in exactly two vertices is covered by the set $C = \{p_1, p_2\}$. Additionally, any edge that intersects $e$ in three vertices and contains the vertex $p_1 \cap p_2 = 1$ is covered by $C$. This implies $\tau(2)(H) \leq \tau(2)(H) \leq 2$ unless there exists an edge containing the vertices $e \setminus (p_1 \cap p_2) = 234$.

---

**THE ELECTRONIC JOURNAL OF COMBINATORICS** 29(4) (2022), #P4.14

---

12
Suppose \( g = 234v, \ v \in V(H) \setminus \{1\}, \) is an edge in \( H. \) It is easy to see that \( \{p_1, p_2, 23\} \) is a cover, implying \( \tau^3(H) \leq 3 \). Notice that \( g \) must intersect both \( f_1 \) and \( f_2 \) in two vertices, implying that \( v \in f_1 \cap f_2. \) If \( u \neq 6 \) then \( f_1 \cap f_2 = 15 \) and, hence, \( g = 2345. \) Observe that the pairs 46 and 4u are not contained in any edge. Indeed, if 46 is in an edge, then \( G[1246] \) and \( G[135u] \) are two disjoint \( K_4 \)'s in \( G. \) If \( 4u \) is in an edge, then \( G[134u] \) and \( G[1256] \) are two disjoint \( K_4 \)'s in \( G. \)

Let \( t : \binom{V(H)}{2} \rightarrow \mathbb{R} \) be defined as follows:

\[
t(x) = \begin{cases} \frac{1}{3} & \text{if } x \in \{12, 13, 15, 23, 25, 35 \} \\ 0 & \text{otherwise} \end{cases}
\]

Then \( t \) is a fractional cover implying \( \tau^*(H) \leq |t| = 2. \) To see that \( t \) is a fractional cover, note that any edge that is a witness for \( p_1 = 12 \) or \( p_2 = 13 \) must intersect \( g \) in at least two vertices and, hence, contains the vertex 5. Otherwise, if an edge intersects \( e \) in three vertices, it must contain 123 or the vertex 5. Indeed, if an edge intersects \( e \) in 124, then in order to intersect \( f_2 \) in at least two vertices, the edge must contain 5 (by the discussion above, it cannot contain the pair 4u). A similar argument shows that an edge that intersects \( e \) in 134 contains the vertex 5.

Now suppose \( u = 6, \) then \( v \in f_1 \cap f_2 = 156. \) If \( H \) contains both 2345 and 2346 as edges, then \( G[123456] = K_6. \) As before, it follows that \( H = \binom{[6]}{4} \) implying \( \tau^*(H) \leq 2.5, \) and \( \tau^2(H) \leq 3. \) We have already dealt with the case where 2345 is an edge. If 2346 is an edge, then we may obtain a fractional cover by replacing the vertex 5 with the vertex 6 in the fractional cover \( t \) defined above. The proof follows similarly.

**Case 2:** There exists an indispensable pair \( p_3 \) such that \( p_1, p_2, p_3 \) form a triangle in \( G, \) i.e., \( p_3 = 23. \) Then \( C = \{p_1, p_2, p_3 \} = \{12, 13, 23 \} \) is a cover. Clearly edges that intersect \( e \) in exactly two vertices are covered. If an edge intersects \( e \) in at least three vertices, then it must contain at least two elements of the set \( p_1 \cup p_2 \cup p_3 = 123 \) and must be covered by \( C. \) It follows that \( \tau^3(H) \leq 3. \)

We now prove that \( \tau^*(H) \leq 2.5. \) Recall that \( e = 1234, \) \( p_1 = 12 \) and \( p_2 = 13 \) have witnesses \( f_1 = 1256 \) and \( f_2 = 135u \) (with \( u \in V(H) \setminus 24 \)). Let \( F_1, F_2 \) and \( F_3 \) be the set of witnesses of \( p_1, p_2 \) and \( p_3 \) respectively. Note that, for any \( i \neq j, \) exactly one of the following holds:

(a) There exist witnesses \( f_i \in F_i \) and \( f_j \in F_j \) with \( |f_i \cap f_j| = 2; \)

(b) \( F_i = \{f_i \}, F_j = \{f_j \} \) and \( |f_i \cap f_j| = 3. \)

It suffices to consider the following cases:

**Case 2.1:** All pairs \( i \neq j \) satisfy (b). Without loss of generality, let \( F_1 = \{f_1 \} = \{1256 \}, F_2 = \{f_2 \} = \{1356 \} \) and \( F_3 = \{f_3 \} = \{2356 \}. \) Let \( t : \binom{V(H)}{2} \rightarrow \mathbb{R} \) be given by

\[
t(x) = \begin{cases} \frac{1}{3} & \text{if } x \in \{12, 23, 13 \} \\ \frac{1}{6} & \text{if } x \in \{14, 15, 16, 24, 25, 26, 34, 35, 36 \} \\ 0 & \text{otherwise} \end{cases}
\]
Then $t$ is a fractional cover with $|t| = 2.5$. Clearly every edge in $F_1 \cup F_2 \cup F_3$ is covered. If $f \cap e = 123$, then it is covered. Otherwise if $f \cap e \in \{124, 134, 234\}$, then $f$ must contain either the vertex 5 or the vertex 6 and so must be covered.

**Case 2.2:** There exists a pair $i \neq j$ satisfying (a). Without loss of generality, suppose $f_1 = 1256 \in F_1$ and $f_2 = 1357 \in F_2$. Now any $f_3 \in F_3$ must satisfy $|f_3 \cap f_1|, |f_3 \cap f_2| \geq 2$ and so either contains the vertex 5 or the pair 67. But if 67 were in some edge of $H$, then $G[1567]$ is a $K_4$ implying $\nu^{(2)}(H) > 1$, a contradiction. It follows that every edge in $F_3$ contains the vertex 5.

Suppose every edge in $F_1 \cup F_2$ also contains the vertex 5. Then we define a fractional cover $t: \binom{V(H)}{2} \to \mathbb{R}$ as follows:

$$t(x) = \begin{cases} \frac{1}{3} & \text{if } x \in \{12, 13, 15, 23, 25, 35\} \\ 0 & \text{otherwise} \end{cases}$$

To see that $t$ is a cover, note that edges in $F_1, F_2$, and $F_3$ are covered by pairs in $\{12, 15, 25\}, \{13, 15, 35\}$, and $\{23, 25, 35\}$, respectively. If $f \cap e = 123$, then clearly it is covered. Otherwise, if $f \cap e \in \{234, 134, 124\}$ then $|f \cap 1256|, |f \cap 1357| \geq 2$ implies that $f$ contains the vertex 5 and, hence, is covered. Since $|t| = 2$, this implies the assertion.

Assume now that there is an edge in $F_1 \cup F_2$ that does not contain the vertex 5. Without loss of generality, suppose $f'_1 \subseteq F_1$ is such an edge. Since $f'_1 \cap (f_2 = 1357) \geq 2$, $f'_1$ contains the vertex 7. But then $f'_1$ cannot contain the vertex 6 (since 67 cannot be in an edge of $H$). By definition $f'_1$ cannot contain the vertices 3 or 4. Therefore, we may assume $f'_1 = 1278$.

For convenience, we recall our assumptions: $e = 1234$, $\{f_1 = 1256, f'_1 = 1278\} \subseteq F_1$, $f_2 = 1357 \in F_2$ and that every edge in $F_3$ contains the triple 235. We can also assume that the pair 67 is not contained in any edge of $H$. For the same reason, the pair 58 cannot be contained in an edge (otherwise 1578 in $H$ implying $\nu^{(2)}(H) > 1$).

Any $f_3 \in F_3$ must be of the form 235$u$. Since $|f_3 \cap f'_1| \geq 2$, $f_3$ must contain one of the vertices 7 or 8. But by the discussion in the preceding paragraph $f_3$ cannot contain 8, implying that $F_3 = \{f_3\} = \{2357\}$. This in turn implies that $f_2 = 1357$ is the unique witness of $p_2$, since $|f_2 \cap f| \geq 2$ for each $f \in \{f_1, f'_1, f_3\}$. Let $t: \binom{V(H)}{2} \to \mathbb{R}$ be given by:

$$t(x) = \begin{cases} \frac{1}{3} & \text{if } x \in \{13, 23\} \\ \frac{2}{3} & \text{if } x \in \{35\} \\ 1 & \text{if } x \in \{12\} \\ 0 & \text{otherwise} \end{cases}$$

Clearly $F_1, F_2$ and $F_3$ are covered. If $f \cap e \in \{123, 124\}$ then it is covered. If $f \cap e \in \{134, 234\}$ then $f$ cannot satisfy $|f \cap f_1| \geq 2$ and $|f \cap f'_1| \geq 2$, hence such an edge cannot exist. It follows that $t$ is a cover, implying $\tau^{(2)}(H) \leq \frac{7}{3}$.

**Case 3:** There exists an indispensable pair $p_3$ such that $p_1, p_2, p_3$ form a $K_{1,3}$ in $G$, i.e., $p_3 = 14$. Any edge that intersects $e$ in exactly two vertices must contain one of the pairs $\{p_1, p_2, p_3\}$. If an edge intersects $e$ in three vertices and contains the vertex $p_1 \cap p_2 \cap p_3 = 1$, then...
then it is covered by some \( p_i \). It follows that \( \tau(2)(H) \leq 3 \), unless there exists an edge intersecting \( e \) in exactly the vertices 234. Suppose there exists \( g \) with \( g \cap e = 234 \). Since \(|g \cap f_1| \geq 2\), \( g \) must contain either the vertex 5 or the vertex 6. Without loss of generality, let \( g = 2345 \). But now any edge that witnesses \( p_1, p_2 \), or \( p_3 \) must also contain the vertex 5 (since it must intersect \( g \) in at least two vertices). It follows that the set \( \{15, 12, 34\} \) is a cover implying \( \nu(3)(H) \leq 3 \).

To bound the size of the optimal fractional cover, as in Case 2, it suffices to consider the following cases.

**Case 3.1:** Let \( F_1 = \{f_1\} = \{1256\} \), \( F_2 = \{f_2\} = \{1356\} \) and \( F_3 = \{f_3\} = \{1456\} \). But then \( G[3456] \) is a \( K_4 \), implying that \( H = (\binom{6}{2}) \).

**Case 3.2:** Suppose \( f_1 = 1256 \in F_1 \) and \( f_2 = 1357 \in F_2 \). As in Case 2, this implies that every edge in \( F_3 \) contains the vertex 5. Observe that the pair 46 is not contained in an edge otherwise \( \{1357, 1246\} \) is a 2-matching in \( H \), a contradiction. Similarly the pair 47 is not contained in an edge, otherwise \( \{1256, 1347\} \) is a 2-matching in \( H \). That is, we may assume \( f_3 = 1458 \in F_3 \).

Now any edge \( f'_1 \in F_1 \) must satisfy \( |f'_1 \cap f_3| \geq 2 \) implying that it must contain the vertex 5 or the vertex 8. But if \( 28 \subset f'_1 \), then \( \{1357, 2458\} \) is a 2-matching in \( H \). It follows that every edge in \( F_1 \) must contain the vertex 5. Similarly every edge in \( F_2 \) must contain the vertex 5, since if \( 38 \) is contained in an edge then \( \{1256, 1378\} \) is a 2-matching in \( H \). Now we define a cover \( t : \binom{V(H)}{2} \to \mathbb{R} \) as follows:

\[
t(x) = \begin{cases} 
1 & \text{if } x \in \{15, 23\} \\
0 & \text{otherwise}
\end{cases}
\]

To see that \( t \) is a cover, we note that all type-1 edges contain the pair 15. If \( f \cap e \in \{123, 124, 134\} \), then it contains the vertex 5, and, hence, the pair 15. Otherwise if \( f \cap e = 234 \), then it is covered by 23. Since \( |t| = 2 \), this implies the assertion. \( \Box \)

### 3.3 Proof of Theorem 8 (c)

To see the lower bound, let \( H \in \mathcal{H}_4 \) be the hypergraph with edge set

\[
\{1234, 1256, 3456, 1367, 2467, 1457, 2357\}.
\]

It is easy to check that \( \nu(2)(H) = 1 \) and \( \tau(2)(H) = \nu(2)(H) = 3.5 \). The hypergraph \( H \) appeared in [1, Proposition 3.8], however, the fractional 2-cover and 2-matching numbers of this hypergraph were not noted. To get some intuition about \( H \), let \( e = 1234 \) and observe that \( \binom{6}{2} \) can be decomposed into three pairs of disjoint elements, specifically \( \{12, 34\}, \{13, 24\}, \{23, 14\} \). For each such pair, e.g., \( \{12, 34\} \) we add edges \( f_1 \) and \( f_2 \) such that \( f_1 \cap e = 12 \) and \( f_2 \cap e = 34 \), respectively, and \( f_1, f_2 \) contain the pair 56 (which is needed to ensure that \( \nu(2)(H) = 1 \)). Similarly, we add two edges for \( \{13, 24\} \), and two edges with for \( \{12, 34\} \) while ensuring that \( \nu(2)(H) = 1 \).

For the upper bound, let \( H \in \mathcal{H}_4 \) be a 4-graph with \( \nu_2(H) = 1 \), and \( e = 1234 \) be an edge of \( H \). Suppose all edges intersect \( e \) in exactly two vertices. We may decompose \( \binom{6}{3} \)
into three pairs of disjoint elements $M_1 = \{12, 34\}$, $M_2 = \{13, 24\}$, and $M_3 = \{23, 14\}$. By Corollary 24, for each $i \in [3]$, if there are witnesses for both pairs in $M_i$, there is a pair $q_i$ not in $e$ contained in these witnesses. If there is a witness for only one pair in $M_i$, let $q_i$ be this pair. Let $t : \left(\binom{V(H)}{2}\right) \rightarrow \mathbb{R}$ be defined as follows:

$$
t(x) = \begin{cases} 
\frac{1}{6} & \text{if } x \in \binom{e}{2} \\
\frac{3}{6} & \text{if } x \in \{q_1, q_2, q_3\} \\
0 & \text{otherwise}
\end{cases}
$$

It is easy to see that $t$ is a fractional cover, implying that $\tau^{\ast(2)}(H) \leq |t| = 3.5$.

From here on, let $f$ be an edge intersecting $e$ in exactly three vertices. Assume without loss of generality that $f = 1235$. Let $p_1 = 12, p_2 = 23, p_3 = 13$ and $p_4 = 45$, and note that $C = \{p_1, p_2, p_3, p_4\}$ is a cover of $H$ of size 4. Indeed, for any edge $g$, $|g \cap e|, |g \cap f| \geq 2$ imply the following:

Either $g$ contains one of the pairs $p_1, p_2, p_3$, or $g$ contains $p_4$. \hspace{1cm} (3)

We may assume that for each $i \in [4]$ there is a nonempty set $E_i$ of edges in $H$ containing $p_i$ and no other element of $C$, otherwise $\tau(H) \leq 3$.

For $1 \leq i \leq 4$, let $f_i \in E_i$. Note that $|f_i \cap e| \geq 2$ implies $f_i$ contains exactly one of the vertices 1, 2, or 3. Without loss of generality, let $f_4 = 3456$. This implies that $E_1 \subseteq \{1246, 1256\}$, since $45 \not\subseteq f_1$ and $|f_1 \cap 3456| \geq 2$. Now if every edge in $E_2 \cup E_3$ contains the vertex 6, then we obtain a fractional cover $t : \left(\binom{V(H)}{2}\right) \rightarrow \mathbb{R}$ given by:

$$
t(x) = \begin{cases} 
1 & \text{if } x \in \{45\} \\
\frac{1}{2} & \text{if } x \in \{12, 13, 23, 26, 36\} \\
0 & \text{otherwise}
\end{cases}
$$

is a cover of size $|t| = 3.5$.

Therefore, we may assume that there exists an edge in $E_2 \cup E_3$ that does not contain the vertex 6. Without loss of generality, let $f_2 \in E_2$ be such an edge. Since $45 \not\subseteq f_2$ and $|f_2 \cap f_4| \geq 2$, $f_2$ must contain exactly one of the vertices 4 or 5. We assume that $f_2 = 2347$; the case when $f_2 = 2357$ can be dealt with by interchanging the roles $e$ and $f$ and relabelling the vertices appropriately. Since any edge $f_1 \in E_1$ satisfies $|f_1 \cap f_2| \geq 2$, this implies that $E_1 = \{1246\}$.

For convenience, we reiterate that we may assume that $H$ contains the edges $e = 1234, f = 1235, f_2 = 2347$ and $f_4 = 3456$, and that $E_1 = \{f_1\} = \{1246\}$. As a consequence, we obtain that every edge $f'_2 \in E_2$ must contain the vertex 4 or the vertex 6 (since $|f'_2 \cap f_1|, |f'_2 \cap f_4| \geq 2$), and every edge $f_3 \in E_3$ must contain the vertex 4 or the pair 67 (since $|f_3 \cap f_1|, |f_3 \cap f_2| \geq 2$).

If every edge in $E_2$ contains the vertex 4, a fractional cover $t : \left(\binom{V(H)}{2}\right) \rightarrow \mathbb{R}$ is given by:

$$
t(x) = \begin{cases} 
1 & \text{if } x \in \{45\} \\
\frac{1}{2} & \text{if } x \in \{12, 13, 16, 23, 34\} \\
0 & \text{otherwise}
\end{cases}
$$
implying $\tau^{*}(2)(H) \leq |t| = 3.5$.

We may now assume that there is $f'_2 \in E_2$ such that $f'_2 = 236w$ with $w \in V(H) \setminus 14$.

**Case 1:** Suppose 1245 is not an edge in $H$. Then $t : \binom{V(H)}{2} \to \mathbb{R}$ given by:

$$t(x) = \begin{cases} 
\frac{1}{2} & \text{if } x \in \{13, 23, 24, 26, 34, 36, 45\} \\
0 & \text{otherwise}
\end{cases}$$

is a cover of size $|t| = 3.5$. Indeed any edge that contains two or more pairs in $C$ is covered (since 1245 $\notin H$). Also, edges in $E_1$ are covered by $\{24, 26\}$, edges in $E_2$ are covered by $\{23, 34, 36\}$, and edges in $E_3$ are covered by $\{13, 34, 36\}$. Every edge in $E_4$ must contain one of the vertices 1, 2 or 3, since it must intersect 1234 in a pair. But it cannot contain 1, since otherwise it does not intersect $f'_2$ in a pair, a contradiction. If it contains 2 or 3 then it is covered by $\{45, 24, 34\}$.

**Case 2:** Suppose $g = 1245$ is an edge in $H$. Since $|f'_2 \cap g| \geq 2$, $f'_2 = 2356$. Then $t : \binom{V(H)}{2} \to \mathbb{R}$ given by:

$$t(x) = \begin{cases} 
1 & \text{if } x \in \{23, 45\} \\
\frac{1}{2} & \text{if } x \in \{12, 13, 16\} \\
0 & \text{otherwise}
\end{cases}$$

is a fractional cover of size $|t| = 3.5$. Indeed any edge that contains at least two elements of $C$ is covered. Any edge that contains $p_2$ or $p_4$ is also covered. Edges in $E_1$ are covered by $\{12, 16\}$. Recall that edges in $E_3$ must contain the vertex 4 or the pair 67. But any $f_3 \in E_3$ must also satisfy $|f_3 \cap 1245| \geq 2$ implying that $f_3$ must contain the vertex 4. Finally, $|f_3 \cap 2356| \geq 2$ implies $E_3 = \{1346\}$ and so is covered by $\{13, 16\}$.

This concludes the proof of the theorem.

**3.4 Proof of Theorem 9**

Let $H \in \mathcal{G}_4$ be a 4-graph, and let $M$ be a maximum 2-matching in $H$. Define the function $t : \binom{V(H)}{2} \to \mathbb{R}$ by $t(x) = 1/2$ if $x \in \{e_2^\alpha : e \in M\}$, and $t(x) = 0$ otherwise. Note that $|t| = 3|M|$ and that $t$ is a fractional 2-cover of all edges in $H \setminus T_1$. It remains to cover type-1 edges, each of which already has weight 1/2 by $t$. To finish the proof, it suffices to show that, for any $e \in M$, edges in $T_1(e)$ can be covered using additional weight at most 1. This follows from the fact, proved ahead, that there exist two pairs such that every edge in $T_1(e)$ contains at least one of them. Indeed, assigning weight 1/2 to both these pairs suffices to cover all edges in $T_1(e)$.

Suppose that there are two disjoint indispensable pairs $p_1, p_2$ with witnesses $f_1, f_2$. By Corollary 24, $f_1 \cap f_2 = q$ is a pair disjoint from $e$ and $H[e \cup q] = \binom{e \cup q}{4}$. In particular, every pair contained in $e$ in indispensable and is witnessed by an edge that contains $q$. Suppose $f \in T_1(e)$ does not contain $q$ and set $p = f \cap e$. Then the edge $f' \in F(e)$ which witnesses $e \setminus p$ and contains $q$ shares at most one vertex with $f$, contradicting Observation 23. It follows that every edge in $T_1(e)$ contains $q$.

Since out of any four indispensable pairs two must be disjoint, from here on, we may assume that there are exactly three indispensable pairs. Without loss of generality, let
$e = 1234$, $p_1 = 12$, $p_2 = 13$ and $p_3$ be indispensable pairs. For $i \in [3]$, denote by $F_i$ the set of edges witnessing $p_i$. We may also assume that, for $i \neq j$, edges in $F_i \cup F_j$ do not contain a common pair $p$, otherwise $p \subseteq \bigcap_{f \in F_i \cup F_j} f$ and all edges in $T_1(e)$ contain either $p$ or $p_k$, where $k \notin \{i, j\}$. In particular, for $i \neq j$, we have $|\bigcap_{f \in F_i \cup F_j} f| = 1$ implying that there are at least three edges in $F_i \cup F_j$ (any two edges in $T_1(e)$ intersect in at least two vertices). Also since, by assumption, both $F_i$ and $F_j$ are non-empty, we have that there exist edges $f_i \in F_i$ and $f_j \in F_j$ such that $|f_i \cap f_j| = 2$.

The above discussion lets us assume, without loss of generality, that $f_1^1, f_1^2 \in F_1$ and $f_2 \in F_2$ with $f_1^1 = 1256$, $f_2 = 1367$ and $|f_1^1 \cap f_2^1 \cap f_2| = 1$. Since $|f_1^2 \cap f_2^2| \geq 2$, $f_2^1$ must contain the vertex $7$, i.e., $f_2^1 = 127u$, $u \in V(H) \setminus V(H)$. We may also assume that $u \neq 5$, since otherwise $G[2567]$ is a $K_4$ implying $2657 \in H$. But $2657$ intersects every edge of $M$ in at most one vertex (a pair contained in an edge of $T_1(e)$ cannot be contained in any other edge in $M$) implying that $M \cup \{2567\}$ is a matching of size greater than $|M|$, a contradiction to the maximality of $M$. From here on, we may assume that $f_2^1 = 1278$.

In the preceding paragraph, we use the following observation about $H$ and the corresponding graph $G$. Since it is used again throughout the proof, we make it explicit. Let $p \subseteq f$ be a pair where $f \in T_1(e)$. Any edge in $M \setminus e$ (a $K_4$ in $G$) can share at most one vertex with $p$ (which corresponds to an edge of $G$). Hence, if there exists an edge $e'$ containing at most one vertex in $e$ such that all pairs in $e'$ are contained in an edge of $T_1(e)$ then $M \cup e'$ is a matching of size greater than $|M|$, a contradiction.

For each $u \in 5678$, the pair $1u$ is contained in $f_1^1$ or $f_2^1$. The pairs 56, 67 and 78 are contained in $f_1^1, f_2$ and $f_2^1$ respectively. By the preceding paragraph, it follows that neither of the pairs 57 or 68 can be contained in an edge of $H$ (otherwise 1567 or 1268 is an edge of $H$). Let $f_3$ be a witness for $p_3$, and note that $|f_3 \cap f| \geq 2$ for each $f \in \{f_1^1, f_2^1, f_2\} = \{1256, 1278, 1367\}$ implies that $f_3$ must contain one of the pairs in $\{57, 67, 68\}$. We have already argued that 57 and 68 cannot be contained in edges of $H$, so $f_3$ is determined uniquely. It suffices to consider the following cases:

**Case 1:** $p_3 = 23$ implying that $f_3 = 2367$. This implies that every edge in $F_2$ contains the pair 36. Indeed, let $f \in F_2$ be an edge that does not contain 36, i.e., $f$ does not contain 6. Since $|f \cap 1256| \geq 2$ and $|f \cap 1278| \geq 2$, $f$ must contain the vertices 5 and 8 (57 cannot be contained in an edge). That is, $f = 1358$. But then $f$ and $f_3 = 2367$ are edges in $T_1(e)$ that share less than two vertices, contradicting Observation 23.

We have shown that every edge in $F_2 \cup F_3$ contains the pair 36, and so every edge in $T_1(e)$ contains at least one of the pairs 12 or 36.

**Case 2:** $p_3 = 14$ implying that $f_3 = 1467$. But then 3467 $\in H$ (since 367 $\in f_2$ and 34 $\in e$) and 1278 are edges in $T_1(e)$ that share less than two vertices, a contradiction. □

4 The function $j^*$

4.1 The hypergraph Turán Problem

Recall that $\text{ex}_r(n, k)$ is the maximum number of edges in an $r$-graph on $n$ vertices that does not contain a copy of $K_r^k$, the complete $r$-graph on $k$ vertices. Let $T(n, k, r)$ be
the minimum number of edges in an \( r \)-graph \( H \) on \( n \) vertices such that any subset of \( k \) vertices contains at least one edge of \( H \). Note that if \( H \) is an \( r \)-graph on \([n]\) such that any subset of \( k \) vertices contains at least one edge of \( H \), and \( G \) is the complement of \( H \), i.e., \( G = \{e \in \binom{[n]}{r} : e \not\in H\} \), then \( G \) is an \( r \)-graph on \([n]\) that does not contain a copy of \( K^*_k \), and vice-versa. It follows that \( T(n, k, r) + \operatorname{ex}_r(n, k) = \binom{n}{r} \). Let

\[
T(n, k, r) = \lim_{n \to \infty} \frac{T(n, k, r)}{\binom{n}{r}}, \quad \text{and} \quad \pi(k, r) = \lim_{n \to \infty} \frac{\operatorname{ex}_r(n, k)}{\binom{n}{r}}.
\]

Then we have \( T(n, k, r) + \pi(k, r) = 1 \). The ratio \( T(n, k, r)/\binom{n}{r} \) is non-decreasing (see e.g. [16]), hence the limits above exist and, for any \( n \),

\[
T(n, k, r) \leq t(k, r) \binom{n}{r}, \quad \text{and} \quad \operatorname{ex}_r(n, k) \geq \pi(k, r) \binom{n}{r}.
\] (4)

Given an \( r \)-graph \( H \), a \( K^*_k \)-cover of \( H \) is a set \( C \) of edges such that every \( K^*_k \) in \( H \) contains at least one edge of \( C \). The number \( T(n, k, r) \) is precisely the size of the smallest \( K^*_k \)-cover of \( H = K^*_n \). Let \( \tilde{t}(r) = \inf\{t \in \mathbb{R} : \forall \text{\( r \)-graphs } H \exists \text{ a } K^*_r \text{-cover of size } \leq t|H|\} \). Clearly

\[
t(r + 1, r) \leq \tilde{t}(r), \quad \text{and} \quad \pi(r + 1, r) \geq 1 - \tilde{t}(r).
\] (5)

We show that some constructions that give bounds on \( t(r + 1, r) \) also give bounds for \( \tilde{t}(r) \). It is well known that \( \tilde{t}(2) \leq 1/2 \). Turán [20] showed that \( t(4, 3) \leq 4/9 \), and conjectured that this was optimal. Based on Turán’s construction, we have the following.

**Lemma 26.** For any 3-graph \( H \), there exists a \( K^*_3 \)-cover of size at most \( \frac{4}{9}|H| \). That is \( \tilde{t}(3) \leq 4/9 \).

**Proof.** Let \( V_0 \cup V_1 \cup V_2 = V(H) \) be a uniform random partition of \( V(H) \), and set \( V_3 = V_0 \).

Specifically, each \( u \in V(H) \) is in \( V_i \), \( i \in [3] \), with probability 1/3 independent of other vertices. Let \( C \) be the collection of edges that have all three vertices in one part, or, for some \( 0 \leq i \leq 2 \), have two vertices in \( V_i \) and one vertex in \( V_{i+1} \). It is easy to check that \( H \setminus C \) is \( K^*_3 \)-free. For an edge \( e \in H \), the probability that \( e \in C \) is 4/9. It follows that \( \mathbb{E}|C| = \frac{4}{9}|H| \), which implies the claim.

Extending the same idea to 4-graphs gives:

**Lemma 27.** For any 4-graph \( H \), there exists a \( K^*_4 \)-cover of size at most \( \frac{3}{8}|H| \). That is \( \tilde{t}(4) \leq 3/8 \).

**Proof.** Since the proof is very similar to that of Lemma 26, we only give a sketch. Let \( V_0 \cup V_1 \cup V_2 \cup V_3 = V(H) \) be a uniform random partition of \( V(H) \), and set \( V_4 = V_0 \) and \( V_5 = V_1 \). Let \( C \) be the collection of edges that satisfy one of the following:

- all four vertices are in one part,
- there is one vertex in each part,
• For \(i \neq j\), \(|e \cap V_i| = 2\) and \(|e \cap V_j| = 2\),
• For some \(0 \leq i \leq 3\), \(|e \cap V_i| = 3\) and \(|e \cap V_{i+1}| = 1\),
• For some \(0 \leq i \leq 3\), \(|e \cap V_i| = 3\) and \(|e \cap V_{i+2}| = 1\).

It is easy to check that \(C\) is a cover and that \(E[|C|] = \frac{3}{8}|H|\) implying the assertion. \(\Box\)

For larger \(r\), we use the following construction of Frankl and Rödl [4]. There are constructions that give better bounds, but this suffices for our purposes.

**Lemma 28.** For any \(l \in \mathbb{N}\) and any \(r\)-graph \(H\), there exists a \(K_{r+1}\)-cover of size at most

\[
\left\lfloor \frac{1}{l} + \left(1 - \frac{1}{l}\right)\right\rfloor |H|.
\]

**Proof.** Let \(A_0, \ldots, A_{l-1}\) be a random partition of \(V(H)\), where each vertex is placed into one of \(A_0, \ldots, A_{l-1}\) with probability \(1/l\) independently of other vertices. For \(B \subseteq V(H)\), we define

\[d(B) = |\{i \in \{0, \ldots, l-1\} : B \cap A_i = \emptyset\}|,
\]

and

\[w(B) = \sum_{i=0}^{l-1} i|B \cap A_i|.
\]

For \(0 \leq j \leq l-1\), let \(C_j\) be the family

\[C_j = \{e \in H : (w(e) + j) \mod l \in \{0, \ldots, d(e)\}\}.
\]

Then for every \(0 \leq j \leq l-1\), the set \(C_j\) covers all copies of \(K_{r+1}\) in \(H\). To see this, let \(U \in \binom{V(H)}{r+1}\) be such that \(H[U] = K_{r+1}\). Since there are \(l - d(U)\) indices such that \(U \cap A_i \neq \emptyset\), at least one such index \(i\) must be in

\[(w(U) + j) \mod l, (w(U) + j - 1) \mod l, \ldots, (w(U) + j - d(U)) \mod l.
\]

Let \(x \in U \cap A_i\) and \(e = U \setminus x\). Now, since \(w(e) \equiv w(U) - i \mod l\) and \(d(e) \geq d(U)\), we have \(0 \leq (w(e) + j) \mod l \leq d(e)\) implying \(e \in C_j\).

Since each \(e \in H\) belongs to exactly \(d(e) + 1\) of the families \(C_0, \ldots, C_{l-1}\),

\[
\sum_{j=0}^{l-1} |C_j| = \sum_{e \in H} (d(e) + 1) = |H| + \sum_{i=0}^{l-1} |A_i|,
\]

where \(A_i = \{e \in H : e \cap A_i = \emptyset\}\). Note that

\[
E[|A_i|] = \sum_{e \in H} P\{e \cap A_i = \emptyset\} = |H| \left(1 - \frac{1}{l}\right)^r,
\]

where \(\frac{1}{l} \leq \frac{1}{r}\).
which gives
\[ E \sum_{j=0}^{l-1} |C_j| = \left[ 1 + l \left( 1 - \frac{1}{l} \right)^r \right] |H|. \]

It follows that there exists a collection of covers \( (C_j : 0 \leq j \leq l - 1) \) whose sizes sum up to the expected value. Since we may pick the smallest of these, there exists a cover of size at most
\[ \left[ \frac{1}{l} + \left( 1 - \frac{1}{l} \right)^r \right] |H|. \]

\[ \Box \]

Setting \( l = (r / \ln r) (1 + o(1)) \) in Lemma 28 gives \( \tilde{t}(r) \leq \frac{\ln r + O(1)}{r} \). For our purposes the following very weak, but exact, bound suffices.

**Corollary 29.** For every \( r \geq 5 \), and any \( r \)-graph \( H \), there exists a \( K_{r+1}^r \)-cover of size at most \( \frac{113}{243} |H| \). That is \( \tilde{t}(r) \leq \frac{113}{243} \approx 0.4650 \).

**Proof.** Setting \( l = 3 \) in Lemma 28 implies that, for every \( r \geq 5 \), there is a \( K_{r+1}^r \)-cover of size at most
\[ \left[ \frac{1}{3} + \left( \frac{2}{3} \right)^r \right] |H|. \]  

For \( r = 5 \), we obtain \( \tilde{t}(r) \leq \frac{113}{243} \). Clearly the function in (6) is decreasing in \( r \), implying the assertion. \( \square \)

Lemmas 26 and 27, and Corollary 29 together imply the following.

**Corollary 30.** For \( r \geq 2 \), \( \tilde{t}(r) \leq 1/2 \).

### 4.2 Proof of Theorem 10

For convenience, we let \( \text{ex} = \text{ex}_m(r, m+1) \) for the rest of this section. Suppose, for contradiction, that the assertion is false and let \( H \in \mathcal{H}_r \) be a minimal counterexample. That is, suppose \( H \) satisfies \( \tau^{(m)}(H) > \text{ex} \nu^{(m)}(H) \), and every subhypergraph \( H' \) of \( H \) satisfies \( \tau^{(m)}(H) \leq \text{ex} \nu^{(m)}(H) \). Let \( g \) be a minimum fractional \( m \)-cover, and \( f \) be a maximum fractional \( m \)-matching respectively.

Let \( U \) be the collection of \( m \)-sets \( u \in \binom{V(H)}{m} \) for which \( g(u) > 0 \). By complementary slackness,
\[ |U| = \sum_{u \in U} 1 = \sum_{u \in U} \sum_{e \in H, e \supset u} f(e) = \binom{r}{m} \nu^{(m)}(H). \]  

First assume that there exists \( u \in U \) with \( g(u) \geq 1 / \text{ex} \), and let \( H' \) be the hypergraph obtained from \( H \) by removing all edges containing \( u \). Then we have
\[ \tau^{(m)}(H) \leq \tau^{(m)}(H') + 1, \quad \text{and} \quad \tau^{(m)}(H) \geq \tau^{(m)}(H') + 1 / \text{ex}. \]
It follows that
\[
\tau^{(m)}(H) \leq \tau^{(m)}(H') + 1
\]
\[
\leq \text{ex} \tau^{*\mu(m)}(H') + 1
\leq \text{ex} \left( \tau^{*\mu(m)}(H) - \frac{1}{\text{ex}} \right) + 1
= \text{ex} \tau^{*\mu(m)}(H),
\]
contradicting the assumption on $H$.

We may now assume $g(u) < 1/\text{ex}$ for each $u \in U$. In particular, every edge $e \in H$ contains at least $\text{ex} + 1$ of the $m$-sets in $U$. This implies that, if we consider $U$ as an $m$-graph on $V(H)$ then every edge of $H$ corresponds to an $r$-set of vertices containing a copy of $K_{m+1}$. By the definition of $\tilde{\ell}(m)$, for every $\epsilon > 0$ there exists a $K_{m+1}$-cover $C$ of $U$ such that $|C| \leq (\tilde{\ell}(m) + \epsilon)|U|$. Note also that every edge of $H$ contains at least one $m$-set in $C$, i.e., $C$ is an $m$-cover for $H$. It follows that
\[
\tau^{(m)}(H) \leq (\tilde{\ell}(m) + \epsilon)|U| = (\tilde{\ell}(m) + \epsilon) \binom{r}{m} \nu^{*\mu(m)}(H)
\leq (1 - \tilde{\ell}(m) + \epsilon) \binom{r}{m} \nu^{*\mu(m)}(H) \leq \text{ex}_m(r, m + 1) \nu^{*\mu(m)}(H) + \epsilon \binom{r}{m} \nu^{*\mu(m)}(H),
\]
where the second to last inequality follows from Corollary 30, and the last inequality follows from (4), (5). Taking $\epsilon \rightarrow 0$, we obtain the result. \hfill \Box

**Acknowledgement**

We are grateful to the referees who made many helpful suggestions to improve the clarity of our presentation.

This paper was written as part of the Iowa State Mathematical Research Teams. We are grateful to the Department of Mathematics at Iowa State University for supporting this project.

**References**

[1] R. Aharoni and S. Zerbib. A generalization of Tuza’s conjecture. *Journal of Graph Theory*, 94(3):445–462, 2020.

[2] J. Baron and J. Kahn. Tuza’s conjecture is asymptotically tight for dense graphs. *Combinatorics, Probability and Computing*, 25(5):645–667, 2016.

[3] F. Chung and L. Lu. An upper bound for the Turán number $t_3(n, 4)$. *Journal of Combinatorial Theory, Series A*, 87(2):381–389, 1999.

[4] P. Frankl and V. Rödl. Lower bounds for Turán’s problem. *Graphs and Combinatorics*, 1(1):213–216, 1985.
[5] Z. Füredi. Maximum degree and fractional matchings in uniform hypergraphs. Combinatorica, 1(2):155–162, 1981.

[6] V. Guruswami and S. Sandeep. Approximate hypergraph vertex cover and generalized Tuza’s conjecture. Proceedings of the 2022 Annual ACM-SIAM Symposium on Discrete Algorithms (SODA), 927–944, 2022.

[7] P. Haxell. Packing and covering triangles in graphs. Discrete mathematics, 195(1-3):251–254, 1999.

[8] P. Haxell, A. Kostochka, and S. Thomassé. A stability theorem on fractional covering of triangles by edges. European Journal of Combinatorics, 33(5):799–806, 2012.

[9] P. Keevash. Hypergraph Turán problems. Surveys in combinatorics, 392:83–140, 2011.

[10] M. Krivelevich. On a conjecture of Tuza about packing and covering of triangles. Discrete Mathematics, 142(1-3):281–286, 1995.

[11] L. Lovász. Minimax theorems for hypergraphs. In Hypergraph Seminar, pages 111–126. Springer, 1974.

[12] W. Mantel. Problem 28. Wiskundige Opgaven, 10(60-61):320, 1907.

[13] K. Markström. Extremal hypergraphs and bounds for the Turán density of the 4-uniform $K_5$. Discrete mathematics, 309(16):5231–5234, 2009.

[14] D. Mubayi, O. Pikhurko, and B. Sudakov. Hypergraph Turán problem: Some open questions. In AIM workshop problem lists, manuscript, 2011.

[15] A. Sidorenko. The method of quadratic forms in a combinatorial problem of Turán. Moscow Univ. Math. Bull, 37:1–5, 1982.

[16] A. Sidorenko. Upper bounds for Turán numbers. Journal of Combinatorial Theory, Series A, 77(1):134–147, 1997.

[17] D. Stinson. Combinatorial designs: constructions and analysis. Springer Science & Business Media, 2007.

[18] M. Szestopalow. Matchings and Covers in Hypergraphs. PhD thesis, University of Waterloo, 2016.

[19] M. Szestopalow. Fractional $K_4$-covers. Graphs and Combinatorics, 33(4):1055–1063, 2017.

[20] P. Turán. On an external problem in graph theory. Mat. Fiz. Lapok, 48:436–452, 1941.

[21] Z. Tuza. A conjecture on triangles of graphs. Graphs and Combinatorics, 6(4):373–380, 1990.