Letter

Quantum state-independent contextuality requires 13 rays

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Abstract

We show that, regardless of the dimension of the Hilbert space, there exists no set of rays revealing state-independent contextuality with less than 13 rays. This implies that the set proposed by Yu and Oh in dimension three (2012 \textit{Phys. Rev. Lett.} \textbf{108} 030402) is actually the minimal set in quantum theory. This contrasts with the case of Kochen–Specker sets, where the smallest set occurs in dimension four.

Keywords: quantum contextuality, Kochen–Specker theorem, graph theory, foundations of quantum physics

1. Introduction

Fifty years ago, Kochen and Specker [1] answered the following question: Is it possible that, independently of which is the quantum state, the quantum observables each possess a definite single value, regardless of whether they are measured or not? The Kochen–Specker (KS) theorem states that this is impossible if the dimension of the underlying Hilbert space is larger than two. One consequence of this theorem is the impossibility of reproducing quantum theory in terms of noncontextual hidden variable theories, defined as those in which the outcomes are independent of the context. A context is a set of mutually compatible quantum observables. In this sense, quantum theory is said to exhibit contextuality.

The original proof of the KS theorem had two other distinctive traits: (i) it only used a finite set of observables with two outcomes, where one outcome is represented by a rank-one projection onto a ray of the Hilbert space. Hereafter, as it is common in the literature, we will...
use ray as synonym of self-adjoint rank-one projection. (ii) The set is KS-uncolorable, i.e., it is impossible to assign values 1 or 0 to each ray while respecting that two orthogonal rays cannot both have assigned 1, and 1 must be assigned to exactly one of \(d\) mutually orthogonal rays. These restrictions are motivated by the observation that orthogonal rays correspond to mutually exclusive outcomes of a sharp observable and \(d\) mutually orthogonal rays correspond to an exhaustive set of mutually exclusive outcomes for a Hilbert space of dimension \(d\).

KS-uncolorable sets of rays are called KS sets [2].

The original KS set had 117 rays in \(d = 3\), which can be grouped in 132 contexts. There have been many efforts for finding simpler sets exhibiting state-independent contextuality (SIC). For instance, Peres and Mermin realized that, by considering observables not represented by rank-one projections and replacing KS uncolorability by a similar condition, one can find very compact sets of observables in \(d = 4\) and \(d = 8\) [3, 4]. Still, these sets can be rewritten in terms of KS sets [5, 6]. So far, it has been shown [2] that the KS set of minimum cardinality occurs in \(d = 4\) and has 18 rays [7]. It also has been proved [2] that, in \(d = 3\), the KS set with minimum cardinality has more than 22 and less than 32 rays [8]. On the other hand, the KS set with minimum number of contexts known occurs in \(d = 6\) and has seven contexts (and 21 rays) [9].

A big step was the observation that SIC based on rays does not need to rely on KS-uncolorable sets. It is enough that they lead to a state-independent violation of a noncontextuality inequality. This substantially simplifies the methods for revealing SIC in \(d = 3\). Specifically, Yu and Oh singled out one set with 13 rays in \(d = 3\) [10]. The optimal state-independent noncontextuality inequalities for this set were identified in [11]. Sets of rays having a state-independent violation of a noncontextuality inequality are called SIC sets.

Recent experiments testing SIC [12–20] and an increasing number of applications, such as device-independent secure communication [21], local contextuality [22, 23], Bell inequalities revealing full nonlocality [24], state-independent quantum dimension witnessing [25], and state-independent hardware certification [19], have stimulated the interest in the following question: Which is the minimal set of rays needed for SIC? It is known that, for \(d = 3\), the answer is 13 [26], but it would be well possible that the minimal set occurs in some higher dimension, as it happens for KS sets. Here we prove that this is not the case.

2. Main result

The basis of our proof is a condition identified by Ramanathan and Horodecki [26, 27] to be necessary for any SIC set in dimension \(d\), namely that the orthogonality graph \(G\) of the set of rays has fractional chromatic number \(\chi_f(G) > d\). The orthogonality graph of a SIC set is the graph in which orthogonal rays are represented by adjacent vertices. A coloring of \(G\) is an assignment of colors to the vertices such that adjacent vertices are associated with different colors. \(\chi_f(G)\) is the infimum of \(\frac{d}{S}\) such that vertices have a set of \(b\) associated colors, out of \(a\) colors, where adjacent vertices have associated disjoint sets of colors.

Instead of considering all possible SIC sets of size \(n\), we rather investigate all graphs with \(n\) vertices. Then, we consider the nondegenerate orthogonal representations (ORs) of any graph \(G\). An OR is an injection \(\phi\), mapping the vertices of \(G\) to rays, such that adjacent vertices in \(G\) are mapped to orthogonal rays. The OR is faithful (FOR) if, conversely, any two orthogonal rays correspond to an edge of \(G\). We denote by \(\Xi(G)\) the smallest dimension of the Hilbert space which still admits a FORs of \(G\). It then follows from the Ramanathan–Horodecki condition that \(G\) is the orthogonality graph of a SIC set only if \(\chi_f(G) > \Xi(G)\). Our main results is then as follows.
Theorem 1. Any graph \( G \) with 12 or less vertices has \( \chi_f(G) \leq \Xi(G) \).

Hence, according to quantum theory, no SIC set with less than 13 rays exists.

3. Proof of theorem 1

We proceed by an exhaustive search for a counterexample, examining all 166 122 463 890 nonisomorphic graphs with up to 12 vertices. Applying a cascade of filters we eventually discard all graphs and prove this way theorem 1. We start by introducing the criteria for defining these filters and then explain our procedure providing intermediate results for each of the filters.

We denote by \( V(G) \) and \( E(G) \) the sets of vertices and edges of \( G \), respectively. The complement \( \bar{G} \) of \( G \) is a graph that has the same vertices while the edges are the complemented set, i.e., \( e \in E(\bar{G}) \) if and only if \( e \notin E(G) \). A subgraph \( S \) of \( G \) is any graph with \( V(S) \subseteq V(G) \) and \( E(S) \subseteq E(G) \). A subgraph is induced if \( S \) is also a subgraph of \( \bar{G} \). It is a simple observation that any \((\text{F})\text{OR}\) is also a \((\text{F})\text{OR}\) of any (induced) subgraph. Defining \( \xi \) analogously to \( \Xi \), but for \( \text{ORs} \), this proves the following.

Lemma 2. By definition, \( \xi(G) \leq \Xi(G) \). If \( S \) is a subgraph of \( G \), then \( \xi(S) \leq \xi(G) \). Similarly, if \( S \) is an induced subgraph of \( G \), then \( \Xi(S) \leq \Xi(G) \).

The union of two graphs \( G_1 \cup G_2 \) consists of the disjoint union of the respective vertex sets and edge sets. The join \( G_1 + G_2 \) of two graphs is the union of both graphs adding one edge between any pair \((v_1, v_2) \in V(G_1) \times V(G_2)\). The graph \( K_1 \) with one vertex and no edge takes a special role in the following simple relations.

Lemma 3. For two graphs \( G_1 \) and \( G_2 \) and \( f \in \{\chi_f, \Xi, \xi\} \), we have \( f(G_1 \cup G_2) = \max\{f(G_1), f(G_2)\} \) and \( f(G_1 + G_2) = f(G_1) + f(G_2) \), with the exceptions \( \Xi(K_1 \cup K_1) = 2 \) and \( \xi(K_1 \cup K_1) = 2 \).

Proof. For \( \chi_f \) the relations are well-known, see, e.g., [29], section 3.10. For \( \Xi \) and \( \xi \) the first relation, the maximum is at least a lower bound, since any \((\text{F})\text{OR}\) of \( G_1 \cup G_2 \) must also be a \((\text{F})\text{OR}\) of \( G_1 \) and of \( G_2 \). Conversely, if at least one of the graphs has more than one vertex then also its \((\text{F})\text{OR}\) has at least dimension two. This \((\text{F})\text{OR}\) can then be transformed by a unitary rotation, such that the image of the \((\text{F})\text{ORs}\) of \( G_1 \) and \( G_2 \) are disjoint and also no rays are orthogonal. Hence one can combine any two \((\text{F})\text{ORs}\) of \( G_1 \) and \( G_2 \) to a \((\text{F})\text{OR}\) in the larger of the dimensions of both \((\text{F})\text{ORs}\). The second relation follows at once, noting that \( \{v_1, v_2\} \in E(G_1 + G_2) \) if and only if either \( v_1 \in V(G_1) \) and \( v_2 \in V(G_2) \), or vice versa, or \( \{v_1, v_2\} \in E(G_1) \), or \( \{v_1, v_2\} \in E(G_2) \). Hence \( \phi \) is a \((\text{F})\text{OR}\) for \( G_1 + G_2 \) if and only if it is a \((\text{F})\text{OR}\) for \( G_1 \) and \( G_2 \), and the spans of \( \phi[V(G_1)] \) and \( \phi[V(G_2)] \) are mutually orthogonal.

These relations are useful for our purposes since they imply that, if a graph or its complement is not connected and \( \chi_f(G) > \Xi(G) \), then this must already be true for a subgraph of \( G \). Hence in our search we only need to consider connected graphs the complement of whose are also connected. Another important consequence of lemma 3 is that

5 The orthogonal rank of a graph is also sometimes denoted by \( \xi \) [28], but there the minimum is taken without the restriction that the \( \text{OR} \) is an injection. This yields slightly different properties.
\[ n_2^2, \text{ where } K_{\ell} \text{ is the completely connected graph with } \ell \text{ vertices} \]

This implies \( \Xi(G) \geq n_2^2 + n_1 \) as soon as \( n_2^2 + n_1 \) is a subgraph of \( G \). A weaker form of this condition is that if \( K_{\ell} \) is a subgraph of \( G \), then \( \Xi(G) \geq \ell \).

As a final ingredient to our proof, we use the seven graphs listed in Table 1. If any of those graphs is an induced subgraph \( S \) of \( G \), then \( \Xi(G) \geq \Xi(S) \) applies. The values of \( \Xi(S) \) are obtained by construction, and due to Lemma 3 it is sufficient to study the five graphs in Figure 1. The construction is similar for all five graphs and we demonstrate the method only.

**Table 1.** List of graphs used for filtering via Lemma 2. The graphs \( \text{Caterpillar}^{n_1, \ldots, n_k} \) are linear graphs of length \( k \), where \( n_v \) leaves are added to vertex \( v \). \( H = \text{Caterpillar}^{2,2} \).

- \( C_{n_1}(1, 2) \)
- \( H + K_2 \)
- \( \text{Caterpillar}^{2,2}_1 \)
- \( \text{Caterpillar}^{2,1,1}_1 \)
- \( C_{n_1}(1, 2, 3) \setminus \{v\} \)
- \( H + K_2 \)

| Graph name | In figure 1 | graph6 | \( \Xi \) | Filter | Remaining |
|------------|-------------|---------|----------|--------|-----------|
| \( H \)    | (a)         | Ebtw    | 5        | (3.1)  | 124 220   |
| \( C_{n_1}(1, 2) \) | (b) | Gbijmo | 5        | (3.2)  | 124 216   |
| \( H + K_2 \) | —          | Fbztw   | 6        | (3.3)  | 4 722     |
| \( \text{Caterpillar}^{2,2}_1 \) | (c) | Fbztw   | 6        | (3.4)  | 569       |
| \( \text{Caterpillar}^{2,1,1}_1 \) | (d) | Fbztw   | 6        | (3.5)  | 400       |
| \( C_{n_1}(1, 2, 3) \setminus \{v\} \) | (e) | Ibgnmmjgj | 6 | (3.6)  | 366       |
| \( H + K_2 \) | —          | Gzznnk  | 7        | (3.7)  | 0         |

\( \xi(n_2^2 + mK_1) = 2n + m \), where \( K_\ell \) is the completely connected graph with \( \ell \) vertices [30, 31]. This implies \( \Xi(G) \geq 2n + m \) as soon as \( n_2^2 + mK_1 \) is a subgraph of \( G \). A weaker form of this condition is that if \( K_\ell \) is a subgraph of \( G \), then \( \Xi(G) \geq \ell \).

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for the most complicated case $\text{Ci}_1(1, 2, 3) \setminus \{v\}$, see figure 1(e). The vertices $\{4, 5, 6, 7\}$ form the induced subgraph $K_4$ and, without loss of generality, we can choose $\phi(4) = (1, 0, 0, 0, 0)$, $\phi(5) = (0, 1, 0, 0, 0)$, $\phi(6) = (0, 0, 1, 0, 0)$, and $\phi(7) = (0, 0, 0, 1, 0)$. Since vertex 3 is adjacent to the vertices $\{4, 5, 6\}$ and not adjacent to vertex 7 or 8, and vertex 7 is adjacent to 8, we have $\phi(3) = (0, 0, 0, a, 1)$ with some $a \neq 0$. By similar arguments, $\phi(2) = (0, 0, b, -1/a^*, 1)$ with $b \neq 0$, and, by symmetry, $\phi(8) = (c, 0, 0, 1)$ and $\phi(9) = (-1/c^*, d, 0, 0, 1)$, with $c, d \neq 0$. Using, that vertex 1 is adjacent to the vertices $\{4, 9, 3, 2\}$, we have $\phi(1) = (0, -1/d^*, -x/b^*, -1/a^*, 1)$ with $x = 1 + 1/|a|^2$, and, by symmetry, $\phi(10) = (-1/c^*, -y/d^*, -1/b^*, 0, 1)$ with $y = 1 + 1/|c|^2$. Eventually, vertex 1 and 10 are adjacent, implying $y/|d|^2 + x/|b|^2 + 1 = 0$, which is a contradiction. However, it is straightforward to find a FOR in dimension 6, proving $\Xi[\text{Ci}_1(1, 2, 3) \setminus \{v\}] = 6$.

For all graphs with less than 13 vertices, we discard those graphs which satisfy at least one of the following filter criteria:

| Filter Criteria |
|-----------------|
| (1) $G$ or $\overline{G}$ is not connected. |
| (2.1) $G$ has subgraph $K_t$, where $\chi_f(G) \leq t$. |
| (2.2) $G$ has subgraph $nK_2 + mK_i$, where $\chi_f(G) \leq 2n + m < \chi_f(G) + 1$ and $m \in \{0, 1\}$. |
| (3.1)–(3.7) $G$ has an induced subgraph $S$ from table 1 with $\chi_f(G) \leq \Xi(S)$. |

For obvious reasons, we fall back to a computer-based proof. We use geng from the software package nauty [32, 33] to generate all nonisomorphic graphs. The fractional chromatic number can be obtained by solving the linear program [29, 34],

\[
\text{maximize: } \sum_{v \in V(G)} x_v \\
\text{subject to: } \sum_{v \in I} x_v \leq 1, \text{ for all } I \text{ of } G \\
x_v \geq 0 \text{ for all } v \in V(G),
\]

where $I$ are independent sets of $G$, i.e., sets of vertices where all vertices are mutually nonadjacent. We find optimal solutions to this program using the software package GLPK [35] and verify the correctness of the solution by applying the strong duality of linear programs, using an accuracy threshold of $\epsilon = 10^{-12}$. We approximate the floating-point value obtained for $\chi_f$ by a rational number with less than $\epsilon$ deviation, while constraining the denominator to be not larger than $nm$, where $n$ is the number vertices of $G$ and $m$ is the number of maximal independent sets. This procedure always succeeds and ensures that the calculation of $\chi_f$ is exact, despite floating-point arithmetic being used in intermediate steps.

We apply all filters (1)–(3.7) consecutively so that each filter reduces the number of candidate graphs. The numbers of graphs remaining after each step are shown in table 2, for filters (1), (2.1), and (2.2), and as a function of the number of vertices of the graph. The list of 566 366 graphs remaining after filter (2.2) is available in graph6-format [36]. For the filters (3.1)–(3.7), we show in table 1 the total number of remaining graphs after each filter. No graph remains after applying all filters, which proves theorem 1.

4. Conclusions

Contextuality is a fundamental feature of quantum observables and can be completely detached from any features of the quantum state of the system. This SIC already occurs for
the most elementary case of observables being sharp and having only two outcomes, one of which is nondegenerate; such observables can be represented by rays in a Hilbert space. Here we have shown that SIC with elementary observables requires at least 13 different observables by performing an exhaustive search over all cases with less observables. The Yu–Oh set is an example of such 13 observables and is already realizable on a three-level quantum system, which is the smallest quantum system allowing for contextuality. This is in contrast to the first instances of SIC, the KS sets, where the smallest set cannot be realized on a three-level system. Therefore, fifty years after the discovery of SIC in quantum theory, we finally have the answer to the question of which is the simplest way to reveal it, i.e., which is the smallest set of elementary observables exhibiting SIC.

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\[
\begin{array}{cccc}
\text{Order} & \text{Graphs} & (1) & (2.1) \\
1 & 1 & 1 & 0 \\
2 & 2 & 0 & 0 \\
3 & 4 & 0 & 0 \\
4 & 11 & 1 & 0 \\
5 & 34 & 8 & 0 \\
6 & 156 & 68 & 0 \\
7 & 1044 & 662 & 28 \\
8 & 12346 & 9888 & 456 \\
9 & 274668 & 247492 & 15954 \\
10 & 12005168 & 11427974 & 957882 \\
11 & 1018997864 & 994403266 & 99869691 \\
12 & 165091172592 & 163028488360 & 19715979447 \\
\end{array}
\]
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