ATC Taskload Inherent to the Geometry of Stochastic 4-D Trajectory Flows with Flight Technical Errors

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Abstract—A method to quantify the probabilistic controller taskload inherent to maintaining aircraft adherence to 4-D trajectories within flow corridors is presented. An Ornstein-Uhlenbeck model of the aircraft motion and a Poisson model of the flow scheduling are introduced along with reasonable numerical values of the model parameters. Analytic expressions are derived for the taskload probability density functions for basic functional elements of the flow structure. Monte Carlo simulations are performed for these basic functional elements and the controller taskload probabilities are exhibited.

I. INTRODUCTION

Gridlock is forecast to occur both in the U.S and Europe if the current air traffic practices are not changed. It is widely believed that a technological shift is necessary to accommodate this growth. Along with this technical evolution, a change in procedures and infrastructure will take place, that will give more autonomy to aircraft and rely on higher-precision, less fault-tolerant operations and trajectories.

The NextGen concept of operations [1] envisions flow corridors, where aircraft follow high-precision 4-D-metered trajectories, across relatively compact regions of airspace. These Required Navigation Performance (RNP) trajectories with increasingly more stringent precision requirements use equipment that is certified with probabilistic tolerance bounds [2]. Airspace precision standards in the U.S. differ widely depending on the phases of flight and airspace blocks, ranging from an accuracy of 0.1 NM (RNP-0.1) for precision approach-to-landing to 10 NM (RNP-10) in areas of the Pacific Ocean [3].

Flow corridors rely on self-separation using Automatic Dependent Surveillance-Broadcast (ADS-B). This equipment allows the construction of high-density, high-altitude en route sectors [4]. The role of controllers is to monitor traffic and intervene in case of extreme - and potentially dangerous - deviations from planned 4-D trajectories.

A unified and provable theory linking traffic complexity, airspace block capacity and controller intervention currently does not exist. This research aims to contribute to the unification of these disparate concepts by introducing tools to clarify the dependency between air traffic flow corridor geometry and the controller taskload needed to maintain the structure. Taskload is defined here as the rate at which the controller must intervene to return aircraft that have reached the tolerance bounds outside of their 4-D trajectories. After an overview of the state of the art, in Section III we introduce fundamental stochastic models of the aircraft flight (an Ornstein-Uhlenbeck process) and of the flow scheduling (a Poisson process), along with reasonable numerical values of the model parameters.

Precisely estimating the parameters of the models is highly relevant, since overestimating the variability may cause inefficiency while underestimating it may cause unsafe conditions and elevated controller taskload. The Ornstein-Uhlenbeck aircraft model is calibrated to match simulated data, obtained from a random number generator defined by a Johnson unbounded system $S_U$. Using simulated data is required in the absence of experimental navigation recordings; the random number generator provides fictitious data consistent with statistical studies of the aircraft motion. The Poisson flow model uses intensity parameters given in recently published research.

In Section IV we then derive analytic expressions for the taskload probability density function for basic functional elements of the flow structure: a single lane, multiple parallel lanes, and flow corridor crossings. In Section V, we show taskload probabilities found by Monte Carlo simulations of these basic functional elements.

II. BACKGROUND

Probabilistic models of air traffic are not new to the literature. This is especially true in the case of conflicts and collisions, where the ICAO has in fact defined the acceptable levels of fatal accident risk at one mid-air collision (physical incrossing) per $10^9$ flight hours [6]. The classic model used to predict risk values for some basic types of air traffic management, known as the Reich collision model, was developed in 1964 [7].

Over the last few years, Bakker and Blom generalized the Reich model and developed collision risk models based on first hitting times of Markov processes, using hybrid-state Markov processes with switching coefficients, or Petri nets [8]. Irvine showed that the minimum relative displacement
between conflicting aircraft can follow a normal distribution, and that the minimum distance between them has a folded normal distribution [9].

Much of the conflict detection research invokes position uncertainty. The classic characterization of position uncertainty was suggested by Paielli and Erzberger and models the error as normally distributed with an time-wise linear standard deviation of the along-track error and a constant standard deviation of the cross-track error [10]. This model has been expanded by Blin et al. to include a dynamic resulting from superposed Wiener processes (Brownian motion) affecting position and velocity, or by switching from a position-based to a time-based random process: the time of arrival at a trajectory waypoint becomes the random variable while the position is deterministic [11].

Stochastic conflict detection research is usually geared at designing resolution algorithms. Significant work in this field has been undertaken by Hu, Prandini et al., who have modeled aircraft oscillations around their prescribed trajectories by means of a Brownian motion, the relative position between two aircraft by time-inhomogeneous Markov chains, introduced spatially correlated wind perturbations, and proposed decentralized algorithms that use potential field methodology or the reachability analysis of a switching diffusion model approximated by a Markov chain [12].

Despite all the aforementioned research, not much has been done in terms of stochastic traffic flow management and route structure design. Moreau et al. [13], or Wan and Roy [14], compared management strategies using Poisson flow models. The Poisson flow model has nevertheless raised some doubts: Schmidt studied conflicts at route intersections and showed the Poisson flow model found elsewhere in the literature for aircraft uncertainty, and is meant to reproduce the behavior of an imperfect flight management system (modeled in the volatility of the Brownian oscillation), and a controlling cockpit (pilot and guidance equipment). The cockpit attempts to prevent major deviations (modeled in the elasticity) from the deterministic trajectory (the mean). To our knowledge, this model has not been previously used in this context.

The controller taskload due to managing individual aircraft can be calculated as the number of interventions needed to prevent an aircraft from trespassing the allowed precision bounds. The authors recognize that this is only one aspect of the controller taskload and of the subsequent workload, and also relies on the assumption that the pilot would not become aware of the deviation and to react (hence changing the aircraft dynamics). The probability density function for the first hitting time of a boundary \( \tau^{(k)} = \inf\{t : X_t \geq k\} \) for an origin at \( X_0 \) and a level \( k \) by an unidimensional Ornstein-Uhlenbeck process has the closed-form solution (2) given by Leblanc et al. [20]

\[
P(\tau^{(k)} \in dt) = f_{\tau}(t) = \frac{k - X_0}{\sqrt{2}\pi\sigma^2} \exp\left(\frac{\kappa}{\sigma^2} \sinh(kt)\right) dt \times \ldots \times \exp\left(\frac{\kappa}{\sigma^2} \left(\frac{X_0}{\sigma^2} - \mu\right)^2 - (k - \mu)^2 + \sigma^2 t - \left(\frac{X_0}{\sigma^2}\right)^2 \coth(kt)\right) \]

(2)

From this, a measure of the taskload required for controlling an individual aircraft is obtained as the probability of the number of interventions over a given period \( T_{\text{max}} \). Assuming all realizations of the aircraft trajectory (after each corrective intervention from the controller) are independent, the (discrete) probability (3) is obtained from autoconvoluting the hitting times density:

\[
P[N^{(T)} \geq n] = P[\tau_1 + \ldots + \tau_n \leq T_{\text{max}}] = \int_0^{T_{\text{max}}} \cdots \int_0^{T_{\text{max}}} f_{\tau}(t) dt \]

(3)

Here the \( f_{\tau}(t) \) is the hitting time density function, \( N^{(T)} \) is the number of corrections over time \( T_{\text{max}} \). \{\tau_1 \cdots \tau_n\} are the successive hitting times, and \( \{n-1\} \) represents \( n - 1 \)
Some literature exists which can help quantify the various probability parameters.

The ICAO Performance-based Navigation (PBN) manual [22] defines Path Definition Error (PDE), Flight Technical Error (FTE), and Navigation System Error (NSE), for which it assumes independent, zero-mean Gaussian distributions. PDE occurs when the path defined in the RNAV system does not correspond to the path expected to be flown over the ground, and can be considered to be negligible; FTE relates to the air crew or autopilot’s ability to follow the defined path or track; NSE refers to the difference between the aircraft’s estimated position and actual position. Further ICAO assumptions are that FTE is an ergodic stochastic process within a given flight control mode but that nothing can be said of the NSE due to sensor errors, relative position from navairds, and inertial errors.

Current ICAO standards for PBN demand that the TSE remain equal to or less than the required navigation accuracy (the RNP level) with 95% probability (i.e. 95% of the flight time), and that the TSE has a probability of less than $10^{-5}$ in exceeding twice the required navigation accuracy. Typically, the $10^{-5}$ requirement provides a greater constraint: for a normally distributed cross-track TSE, this bounds the standard deviation to be $\sigma \leq 0.45\text{-}RNP$, while the 95% requirement sets a bound at $\sigma \leq 0.51\text{-}RNP$. In addition, onboard equipment is designed to issue an alert if there is a greater than $10^{-7}$ probability per hour that the error exceed the RNP level by two times.

Numerical and statistical parameters for FTE and FMS performance can be found in the literature, mostly dealing with small aircraft and general aviation [22], [23]. But perhaps the most conclusive statistical analysis was conducted by Levy et al. [24], who identified a Johnson $S_U$ (“unbounded system” [25]) curve as the best fit probability density function for the lateral FTE of a Boeing 747, under RNP-0.3 sensitivity. The fitted curve in [2] is a transformation of the standard normal distribution, and is adapted to heavy-tailed skewed data. An estimate of the curve parameters was conducted by Levy et al. according to a dedicated algorithm developed by Hill et al. in 1976 [26] which matches the first four moments. For the lateral FTE data, the first four moments found by Levy et al. (mean, variance, relative skewness, relative kurtosis) are shown in Table I. The vertical mean and variance are also given by Levy et al., and since no performance data exists for longitudinal data, these were postulated for the purpose of this research by scaling the distribution. If $X_{SNV}$ is a standard normal variate then its transformed $X_{SU}$ obeys [3] and its density function is [9].
adhering to the moments in Table II and the sampling time step for this data was assumed to be 1 minute.

Fig. 2. Simulated FTE data

Fig. 3. Simulated FTE histograms and theoretical density functions

By applying Itō’s lemma, it can be shown that for any fixed $s$ and $t$, $0 \leq s \leq t$, the random variable $X_t$ conditional upon $X_s$ of the Ornstein-Uhlenbeck process has the form (10), where $\mathcal{N}(0, 1)$ denotes a standard normal distribution. The relationship (11) between consecutive observations $X_i$ and $X_{i+1}$ with a timestep $\delta t$ is therefore affine with an independent
and identically distributed random noise $\epsilon$.

$$X_t = X_0 e^{-\kappa (t-s)} + \mu (1 - e^{-\kappa (t-s)}) \ldots$$

$$X_{i+1} = aX_i + b + \epsilon$$  \hspace{1cm} (11)

A least squares linear regression is used to identify the recursion parameters $a$, $b$, and the standard deviation of the noise $a$, from which the parameters (12) of the Ornstein-Uhlenbeck stochastic differential equation (1) can be deduced.

$$\kappa = - \frac{\ln a}{\delta t}$$

$$\mu = \frac{b}{1 - a}$$

$$\sigma = a e^{\kappa \delta t} \left[ \frac{-2 \ln a}{\sigma (1 - a^2)} \right]^{\frac{1}{2}}$$  \hspace{1cm} (12)

A maximum likelihood estimate method was also conducted. From (11), $\{X_{i+1} - aX_i - b = \epsilon\}$ is a normal random variable, and so the conditional probability density function of $X_{i+1}$ given $X_i$ with a time step $\delta t$ is shown in (13), while the log-likelihood function of $n+1$ observations $\{X_0, \ldots, X_n\}$ is given in (14).

$$f_{X_{i+1}|X_i}(x) = \frac{1}{\sqrt{2\pi \sigma^2}} \ldots$$

$$\ldots \times \exp \left[ \frac{(x - X_i e^{-\kappa \delta t} - \mu (1 - e^{-\kappa \delta t}))^2}{2\sigma^2} \right]$$  \hspace{1cm} (13)

$$\hat{\mu} = \sigma \left[ \frac{-2 \ln a}{\sigma (1 - a^2)} \right]^{\frac{1}{2}}$$

$$\mathcal{L}(\mu, \kappa, \hat{\sigma}) = \frac{1}{2\pi \sigma^2} \sum_{i=0}^{n-1} (X_{i+1} - X_i e^{-\kappa \delta t} - \mu (1 - e^{-\kappa \delta t}))^2$$  \hspace{1cm} (14)

The argument of the maximum of $\mathcal{L}$ found from the three partial derivatives $\frac{\partial \mathcal{L}}{\partial \kappa}$, $\frac{\partial \mathcal{L}}{\partial \mu}$, $\frac{\partial \mathcal{L}}{\partial \sigma}$ gives the system in (15).

$$\kappa = - \frac{1}{\delta t} \ln \frac{\sum_{i=0}^{n-1} (X_{i+1} - \mu) (X_i - \mu)}{\sum_{i=0}^{n-1} (X_{i+1} - \mu)^2}$$

$$\mu = \frac{\sum_{i=0}^{n-1} X_{i+1} - X_i e^{-\kappa \delta t}}{n (1 - e^{-\kappa \delta t})}$$

$$\sigma^2 = \frac{1}{n} \sum_{i=0}^{n-1} [X_{i+1} - \mu e^{-\kappa \delta t} (X_i - \mu)]^2$$  \hspace{1cm} (15)

The resulting parameters $\kappa$, $\mu$, $\sigma$ of the Ornstein-Uhlenbeck model - for spatial units expressed in nautical miles (lateral and longitudinal) or feet (vertical) and temporal units expressed in minutes - are shown in Table IV. These values coincide for both the maximum likelihood and least squares estimation methods.

### B. Flow

We adopt a widely accepted model of interarrival times distributed according to an invariant Poisson process; this is equivalent to modeling passage times along the flow by an exponential distribution, and is a fair approximation of the system behavior over a short-time horizon (hourly time frame). If $N(t)$ is a cumulative count of all aircraft having entered the sector, then the probability distribution for the number of aircraft entering a flow between times $t$ and $t + \tau$ is given by (16) and the probability density for the values of interarrival times $x$ is given by (17).

$$P[N(t + \tau) - N(t) = k] = \frac{e^{-\lambda \tau} (\lambda \tau)^k}{k!}$$  \hspace{1cm} (16)

$$f_\lambda(x) = \lambda \exp^{-\lambda x}, x \geq 0$$  \hspace{1cm} (17)

Although most rigorously, the distribution of velocities inside a flow is also a random variable [28], we will assume the uniformity of velocities at 480 knots. This can be justified in the context of 4-D trajectories by constraints added to the along-track performance of the aircraft, which is not the case in current procedures. For a more realistic case, the flow would need to be modeled as a birth-death process with two different Poisson processes, on entry and on exit from the controlled region. This adds some difficulty to the problem, without making it prohibitive.

The statistical study of the typical traffic entering Cleveland center conducted by Salatun et al. [16] provides mean aircraft spacing between 50 and 200 NM for the ten busiest flows in the sector, and an hourly Poisson intensity parameter $\lambda_{center}$ between 10 and 100 (aircraft per hour) depending on time of day, with values averaging 80 for most of the busy times (6 AM to 8 PM EST). For the ten busiest flows in the center, the velocity and spacing values provide average intensity parameters $\lambda_{flow}$ from 10 to 2.5 (aircraft per hour). A reasonable estimate for the time an aircraft spends in a sector is 20 minutes.

From a spatial perspective, current en route flows have wide boundaries, reaching up to 40 NM in width, or even defying any reasonably defined geographic structure in some

| Lateral | Vertical | Longitudinal |
|---------|----------|--------------|
| $\kappa$ | 3.492 min$^{-1}$ | 1.841 min$^{-1}$ | 2.1662 min$^{-1}$ |
| $\mu$ | 2.79 $\cdot$ 10$^{-2}$ NM | 8.034 ft | 9.965 $\cdot$ 10$^{-2}$ NM |
| $\sigma$ | 7.27 $\cdot$ 10$^{-2}$ NM | 8.683 ft | 0.2774 NM |
TABLE V
RNP CORRIDOR MODEL PARAMETERS

| Parameter | Value | Parameter | Value |
|-----------|-------|-----------|-------|
| Intensity (a.c./h) | $\lambda_{inf} = 2.5$ | $\lambda_{sup} = 10$ | |
| Lateral bounds (NM) | $y_{inf} = 0.6$ | $y_{sup} = 1$ | |
| Vertical bounds (ft) | $z_{inf} = 60$ | $z_{sup} = 120$ | |
| Longitudinal bounds (NM) | $x_{inf} = 1$ | $x_{sup} = 5$ | |

cases; we propose that the newly designed flow corridors will concentrate traffic to much tighter regions, with corridor bounds ranging between 3 NM and 10 NM lateral, and 1000-2000 ft vertical, and also supporting parallel lanes, or stacked corridors (see Figure 4). The flow model parameters are shown in Table V.

![Fig. 4. Multilane flow corridors](image)

IV. EXPRESSION OF TASKLOAD PROBABILITIES

In this section we consider several basic functional elements of the flow corridor structure.

A. Single lane

In the case of a single lane of traffic, the overall controller taskload is a superposition of the individual effects from each aircraft. From the Poisson process flow model in (16), the probability of the number of aircraft $M$ simultaneously controlled (a random variable) can be deduced: if it takes an aircraft $T_{cross}$ to cross the airspace, then the probability that at any given time there are $k$ aircraft present is given in (18); if there are $M = k \geq 1$ aircraft present, and the aircraft $i$ requires $N_i(T)$ interventions over time $T$, then the total taskload the controller is subject to has a probability given in (19).

$$P[M = k] = e^{-\lambda T_{cross}} \frac{(\lambda T_{cross})^k}{k!} \quad (18)$$

$$P[N_1(T) + \cdots + N_k(T) = n] = \left[ \sum_{i=1}^{k-1} P\{N(T)\}[^i] \right] P\{N(T)\}[n] \quad (19)$$

Here $\left[ \sum_{i=1}^{k-1} P\{N(T)\}[^i] \right] P\{N(T)\}[n]$ represents $k-1$ discrete autoconvolutions of the single aircraft taskload probability $P\{N(T)\}$ from (4). Obviously for $M = 0$, there will be no intervention with probability 1.

Since the number of aircraft is a random variable, combining (18) and (19) gives the probability (20) of the overall taskload $N_\lambda(T)$ for the flow with intensity $\lambda$ over time $T$ for $n \geq 1$; $P[N(T) = 0]$ is the probability that no intervention is required for one aircraft, given by (5).

$$P[N_\lambda(T) = n] = \sum_{i=1}^{+\infty} P[M = i] \cdot \left[ \sum_{j=1}^{i-1} P\{N(T)\}[^j] \right] P\{N(T)\}[n] \quad (20)$$

B. Multiple parallel lanes

Let us consider the case of parallel lanes (see Figure 4), with either opposing or matching traffic orientation. If the spatial extents of the flows are different, the problem is identical, and assuming independence of the flows, the whole system is equivalent to a single Poisson process with cumulative intensities $\lambda_{tot} = \sum \lambda_i$. Results from Section IV-A can thus be simply generalized to (21).

$$P\{N_\lambda(T) \cup \lambda_i \} = P\{N(T) \lambda_i \} \quad (21)$$

If the spatial extents of the flows are different, the problem is slightly more complex. The taskload probability (22) for $j \geq 2$ different flows with $\{e_1, \cdots, e_j\}$ spatial extents, $\{\lambda_1, \cdots, \lambda_j\}$ intensities, and $\{N_{\lambda_1,e_1}(T), \cdots, N_{\lambda_j,e_j}(T)\}$ interventions per flow is

$$P[N_{\lambda_1,e_1}(T) + \cdots + N_{\lambda_j,e_j}(T) = n] = \left[ \sum_{i=1}^{j} P\{N_{\lambda_1,e_1}(T) \cup \cdots \cup N_{\lambda_j,e_j}(T) \}[^i] \right] P\{N(T)\}[n] \quad (22)$$

where $\left[ \sum_{i=1}^{j} P\{N_{\lambda_1,e_1}(T) \cup \cdots \cup N_{\lambda_j,e_j}(T) \}[^i] \right]$ designates the successive convolutions of the $P\{N_{\lambda_1,e_1}(T) \cup \cdots \cup N_{\lambda_j,e_j}(T) \}$ taskload probabilities given in (20) for flows 1 to $j$. This operator is well defined since convolution is associative.

C. Crossings and mergings

In the case of a crossing, an additional source of taskload comes from possible scheduling conflicts. For two flows $\{F_1\}$ and $\{F_2\}$ crossing at an angle $\alpha$ (in the plane defined by the centerlines of these two flows, see Figure 5) with lateral extents $e_1$ and $e_2$, Poisson intensities $\lambda_1$ and $\lambda_2$, a safe-zone can be defined around the intersection of the centerlines, where only one aircraft is allowed simultaneously. For a minimum possible approach distance $D_{min}$ taken to be for example
5 NM, i.e. the conflict separation standard), the symmetrical safe-zone boundaries $x_1$ and $x_2$ in each of the flows are defined by (23). Figure 6 shows plots of some of these boundary values.

![Fig. 5. Crossing of flows](image1)

![Fig. 6. Possible safe-zone bounds solutions](image2)

\[
\begin{align*}
(x_1 + x_2 \cos \alpha - \frac{e_2}{2} \cos \alpha) & \ldots \\
& \ldots + \left( -\frac{e_1}{2} + x_2 \sin \alpha + \frac{e_2}{2} \sin \alpha \right)^2 = D_{\text{min}}^2 \\
(x_1 + x_2 \cos \alpha - \frac{e_2}{2} \sin \alpha) & \ldots \\
& \ldots + \left( -\frac{e_1}{2} + x_2 \sin \alpha + \frac{e_2}{2} \cos \alpha \right)^2 = D_{\text{min}}^2 \quad (23)
\end{align*}
\]

The safe-zone can then be extended for simplicity so that its crossing time is equal along both flows. The probability of scheduling conflicts is once again simply found from superposing the Poisson processes: if passing the safe-zone of scheduling conflicts is once again simply found from boundary values.

\[
\mathcal{P}[N_{\lambda_1 \cup \lambda_2}^{(T)} = n] = \mathcal{P}[A - 1 + N_{\lambda_1 \cup \lambda_2}^{(T)} = n] \\
= \cdots \mathcal{P}[A - 1] \cdot \mathcal{P}[N_{\lambda_1 \cup \lambda_2}^{(T)}] \\
= \sum_{i=0}^{n} \mathcal{P}[A = i + 1] \cdot \mathcal{P}[N_{\lambda_1 \cup \lambda_2}^{(T)} = n - i] \\
= \mathcal{P}[N_{\lambda_1 \cup \lambda_2}^{(T)} = 0] = \mathcal{P}[A = 0] + \mathcal{P}[A = 1] \cdot \mathcal{P}[N_{\lambda_1 \cup \lambda_2}^{(T)} = 0] \quad (25)
\]

V. MONTE CARLO SIMULATION

Several Monte Carlo simulations were conducted in order to estimate the resulting taskload. Four control tolerance standards (stringent, severe, intermediate, lax) were defined, shown in Table VI. The control bounds are extremely restrictive in order to illustrate the low probabilities of oscillation. The occurrence of a lateral FTE of 0.3 NM has a probability of $5 \cdot 10^{-6}$ in the Ornstein-Uhlenbeck model used in the simulation, which corresponds to a probability of $6 \cdot 10^{-4}$ for one intervention required over 2 hours. For a lateral FTE of 0.4 NM, the probability drops below $2 \cdot 10^{-7}$, which is the resolution limit of the Monte Carlo simulation. For each run, the taskload calculation horizon was $T = 2$ hours.

A. Single lane

Several flow intensities were chosen and a varying number of runs were performed (detailed in Table VII) in order to simulate $2.5 \cdot 10^5$ aircraft for each flow intensity. The lateral containment taskload probability for a two hour interval is shown in Figure 7 and the cumulated lateral, vertical, and longitudinal taskload (3-D control) appears in Figure 8. As expected, the plots show increasing taskload with higher density flows and with stricter tolerance standards. Nevertheless, the values remain relatively low, with no more than 9 or 10 interventions being expected over a 2 hour time frame in all cases.

| TABLE VI | TOLERANCE STANDARDS |
|----------|----------------------|
| Lateral (NM) | Vertical (ft) | Longitudinal (NM) |
| Stringent | 0.1 | 20 | 0.5 |
| Severe | 0.12 | 22 | 0.6 |
| Intermediate | 0.15 | 25 | 0.8 |
| Lax | 0.2 | 30 | 1 |

| TABLE VII | MONOLANE MC SIMULATION PARAMETERS |
|-----------|-----------------------------------|
| Intensity $\lambda$ (a.c/h) | 2.5 | 5 | 7.5 | 10 | 60 |
| MC runs | 91,658 | 66,680 | 58,366 | 54,147 | 41,702 |
| Simulated aircraft | 250,000 |

\[
\mathcal{P}[A = c] = \frac{e^{-\lambda (1+\lambda_2)^{T_{\text{safe}}}}[(\lambda_1 + \lambda_2)T_{\text{safe}}]^c}{c!} \quad (24)
\]

This can also be interpreted as a lower bound of the number of required conflict resolution interventions: $c - 1$ of those aircraft will need to be delayed and/or deviated, without accounting for possible secondary conflicts. This probability however does account for scheduling conflicts within each separate flow, as well as the crossing flows combined.

From (20), (21), and (24), a discrete convolution [$*$] in (25) gives the total taskload over time $T$ for $n \geq 1$ accounting for $A - 1$ conflicts at the flow crossing and for $N_{\lambda_1 \cup \lambda_2}^{(T)}$ interventions to maintain the structure.
The Monte Carlo simulation parameters of the crossing are shown in Table IX. The taskload calculations were performed for two crossing flow intensities of $\lambda = 2.5$ a.c/h. It was assumed that in the enroute domain, the higher density flow corridors would be designed for the express purpose of avoiding such crossings. Simulations show that in the higher density levels, the taskload comes almost exclusively from resolving scheduling conflicts at the crossing, rather than from controlling the aircraft deviations.

Figure 10 shows the effect of the crossing angle $\alpha$ on the controller taskload for a crossing of two flows with stringent deviation tolerance. Since the $\alpha = 90^\circ$ crossing angle imposes the smallest safe-zone extent, it is intuitive that this also corresponds to the lowest conflict resolution taskload, and inversely for the $\alpha = 120^\circ$ crossing which causes the largest safe-zone. The spatial extent also explains the control taskload: the longer the time an aircraft spends within the safe-zone, the higher the probability that it will touch one of the control bounds. Therefore, the $\alpha = 90^\circ$ crossing has the highest probability that no control intervention will be needed, and the $\alpha = 120^\circ$ crossing has the highest probabilities for any of the superior taskload values.

Figure 11 shows the angle effect on a crossing between a stringent and a severe flow. When compared to Figure 10, the taskload probability decreases slightly (best seen for the $\alpha = 90^\circ$ crossing) because of the higher control tolerance in the second flow.

Figures 12, 13 and 14 show the effect of the control tolerance on the taskload at a crossing of $\alpha = 30^\circ$, $\alpha = 90^\circ$, and $\alpha = 120^\circ$ respectively. As can be expected, for all the crossing values the control tolerance bounds play a significant role in the taskload levels. This effect is however least noted in the $\alpha = 90^\circ$ crossing where the aircraft spend the least

### TABLE IX

**Crossing MC simulation parameters**

| Crossing angle $\alpha$ (deg) | 30 | 90 | 120 |
|------------------------------|----|----|-----|
| Crossing time (min)          | 2  | 1  | 3   |
| MC runs                      | 5,412 | 10,463 | 3,826 |
| Simulated aircraft           | 25,000 | 50,000 | 16,666 |

---

**TABLE VIII**

**Multilane MC simulation parameters**

| Intensity $\lambda$ (a.c/h) | 2.5 | 5 | 7.5 | 10 | 60 |
|-----------------------------|-----|---|-----|----|----|
| MC runs                     | 22,9158 | 16,670 | 14,592 | 13,537 | 10,426 |
| Simulated aircraft          | 62,500 |

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**B. Multiple parallel lanes**

As noted in Section IV-B, the case of several flows with identical spatial extents (or control tolerances) and varying Poisson intensities can be simply reduced to a single flow with cumulative intensity. The case of lanes with varying precision requirements poses more interest. In this simulation, the very high density flow intensity parameter (60 a.c/h) was chosen for all the lanes, while the control tolerance standards were stringent for the first flow, severe for the second, intermediate for the third, and lax for the fourth.

The Monte Carlo simulation parameters of the multilane flow are shown in Table VIII. The taskload probability for a two hour interval is shown in Figure 9. The simulation reveals that although taskload increases with additional lanes, the stringent flow is the source of most of the taskload. Adding the severe flow to this has significant impact by increasing the most probable taskload level from 3 to 4 interventions over a two hour time frame in the 3-D control case. Additional flows with more tolerant bounds do not have notable effect, only increasing taskload probabilities by less than 1%.

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**Fig. 7.** Effect of flow rate and control tolerance on monolane lateral control 2h taskload probability

**Fig. 8.** Effect of flow rate and control tolerance on monolane 3-D (lateral, vertical, longitudinal) control 2h taskload probability

**Fig. 9.** Effect of number of lanes on multiline 2h taskload probability
amount of time in the safe-zone. For the intermediate and lax tolerances, the taskload only comes from the conflict resolution interventions (the probability of no control interventions is close to 1). Since the control tolerance plays no part in the flow scheduling, the conflict resolution taskload has not been shown in plots other than in Figure 10.

![Fig. 10. Effect of crossing angle on 2h taskload probability; stringent standards for both flows](image)

(a) Deviation control taskload probability  
(b) Conflict resolution taskload probability  
(c) Total taskload probability

Fig. 10. Effect of crossing angle on 2h taskload probability; stringent standards for both flows

![Fig. 11. Effect of crossing angle on 2h taskload probability; stringent and severe standard flows intersect](image)

(a) Deviation control taskload probability  
(b) Total taskload probability

VI. CONCLUSION

In a context of NextGen flow corridors, this research seeks to quantify what controller taskload is required to maintain aircraft adherence to their 4-D trajectories, within certain tolerance levels. Fundamental stochastic models of the aircraft motion (an Ornstein-Uhlenbeck process) and of the flow scheduling (a Poisson process) have been introduced, along with reasonable numerical values of the model parameters. The Ornstein-Uhlenbeck aircraft model was calibrated - by a least squares linear regression and by a maximum likelihood estimation - to match simulated data, obtained from a random number generator defined by a Johnson unbounded system $S_U$. The use of simulated data to calibrate the model was imposed by the absence of experimental Flight Technical Error (FTE) data; the random number generator provided fictitious data consistent with statistical studies of the aircraft motion. The Poisson flow model used intensity parameters given in recently published research. Analytic expressions have been derived for the taskload probability density functions for basic functional elements of the flow structure (a single lane, multiple parallel lanes, and flow corridor crossings). Monte Carlo simulations have been performed for these basic functional elements of the flow structure and the controller taskload probabilities have been exhibited.

It has been found that the current aircraft performance de facto supports the flow corridor concept, by providing high-precision navigation. The controller taskload needed -
assuming no anticipated reaction from the pilot - to prevent the aircraft from excessive deviation outside of their nominal 4-D trajectories was moderate, not exceeding ten interventions over two hours for a single or multiple lane flow with very high density traffic (a mean of 60 aircraft per hour). The crossing of flows however poses different concerns, where solving the occurring scheduling conflicts at the crossing point becomes the foremost source of controller taskload.

The authors acknowledge that experimental FTE data measurements and a more accurate statistical flow description are needed to improve the model accuracy and, in turn, to increase confidence in its predictions. Nevertheless, a coherent deductive taskload model that only requires simple fundamental microscopic (aircraft motion) and macroscopic (flow) models has been demonstrated in this research.

In future work, we plan on using the detailed Ornstein-Uhlenbeck model of the aircraft developed here to quantify risk severity in the case of extreme deviations from nominal operations. Studies of the risk propagation across flow geometry under degraded operations, and its dependency with the route structure will also be conducted.

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