In praise of measurement

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Abstract. The role of measurement in quantum computation is examined in the light of John Bell’s critique of the how the term “measurement” is used in quantum mechanics. I argue that within the field of quantum computer science the concept of measurement is precisely defined, unproblematic, and forms the foundation of the entire subject.

Here are some words which . . . have no place in a formulation with any pretension to physical precision: system, apparatus, environment, microscopic, macroscopic, reversible, irreversible, observable, information, measurement. On this list of bad words the worst of all is “measurement”. . . . What exactly qualifies some physical systems to play the role of “measurer”? . . . The word has had such a damaging effect on the discussion, that I think it should now be banned altogether in quantum mechanics.

— J. S. Bell [1]

In our description of nature the purpose is not to disclose the real essence of the phenomena but only to track down, so far as it is possible, relations between the manifold aspects of our experience.

— Niels Bohr [2]

In his his elegant tirade against measurement John Bell declared that we lack “an exact formulation of some serious part of quantum mechanics.” He explained that by “exact” he meant “fully formulated in mathematical terms, with nothing left to the discretion of the theoretical physicist.” And by “serious” he meant that “some substantial fragment of physics should be covered” and that “‘apparatus’ should not be separated off from the rest of the world into black boxes.”

It’s sad that Bell’s early death deprived him of the pleasure of experiencing the quantum-computation revolution, and a misfortune for physics that it deprived us of the critical insights he surely would have had into its implications for quantum foundations. I don’t think Bell would have denied that the theory of quantum computation is as exact as
any branch of applied mathematics can be said to be, with nothing left to the discretion of the user beyond the choice of problem to which to apply it, and his or her ingenuity in devising quantum algorithms.

Whether he would have viewed quantum computation as a “serious part of quantum mechanics” is less clear. I would maintain that any fragment of physics, large enough to be applied to the efficient factoring of enormous integers, has to be viewed as substantial. The U.S. National Security Agency surely does. But there is no denying that the computational apparatus — the actual machinery of the computer — is separated off into black boxes from the Qbits on which it acts. [I digress to commend the term “Qbit”, as a highly convenient abbreviation for “qubit”, especially in contexts where one also talks about Cbits, as the physical carriers of a bit of information in a classical computer.] But that machinery is also, unproblematically, separated off into black boxes from the Cbits on which it acts in classical computer science.

I’m not sure Bell would have found any of the remarks that follow compelling, or even suggestive. Nevertheless, I have found it illuminating to reexamine the role of measurement in quantum computation from the perspective of his critique, and I offer such a reexamination as a 60th birthday present to Anton Zeilinger, who I hope will take a more sympathetic view of it than John Bell might have done, if we had had the good fortune to have him with us here today.

The first striking change in how quantum mechanics appears, when viewed through its serious subset of quantum computation, is that the continuous time evolution of ordinary quantum mechanics is replaced by the actions of a collection of discrete gates. Some important quantum-computational gates are shown in Figure 1, together with the gate of special interest to us, the 1-Qbit measurement gate. All of the gates in Figure 1, including the measurement gate, alter the state associated with the incoming Qbits in a well-defined, generally discontinuous manner, which is precisely defined by that state, with nothing left to the discretion of the theoretical physicist. Within this framework, the action of the measurement gate differs from the other gates, all of which are unitary, in several ways.

Most importantly, unlike the unitary gates, the measurement gate gives the user information. Although “information” is another of Bell’s bad words, there is nothing vague about its use here: the number $x$ showing in the user-readable display is either 0 or 1. “Information” can’t get more precise or elementary.

Next in importance, unlike the other gates, a measurement gate has an output characterized by a state that is only statistically determined by the state associated with its input, the probability $p(x)$ being given by the modulus squared of the amplitude $a_x$ characterizing the input state.

The measurement gate also differs from the others in having no inverse (a benign example of Bell’s bad word irreversibility). And its conversion of incoming to outgoing states is not unitary, and, indeed, not even linear.

On the other hand the measurement gate resembles the unitary gates in that the rule
governing the state characterizing its output is completely and precisely specified by the state characterizing its input. Another similarity is that the action of the measurement gate, like that of the unitary gates, is discrete, in contrast to the notorious distinction traditionally made between the continuous evolution of the state under the Schrödinger equation and the discontinuous change of state in a measurement. In most forms of quantum computation the role of the Schrödinger equation is performed by a small number of discrete unitary gates. So time evolution and measurement are alike in both producing discrete changes of state.

The virtues of the 1-Qbit measurement gate are many, whence my title. Most obviously, without measurement gates, a computation has no output. There is in principle
no way, other than reading the display on measurement gates, for the user of a quantum computer to extract information from Qbits. This is less distressing in computer science than it may be in physics, since there is also no way in practice for the ordinary user of a classical computer to extract information from the Cbits, hardwired on microchips buried in the innards of his or her machine, other than by looking at a visual display or a print-out.

A less obvious — or at least less often emphasized — virtue of the 1-Qbit measurement gate is that, at least in the most straightforward formulations of quantum computer science, without measurement gates the computation also has no useful input, as I shall expand upon below. A not unrelated virtue, of great importance for the practical feasibility of quantum computation, is that without measurement gates there can be no error correction.

A final virtue, relevant to John Bell’s critique of measurement, is that the many complex, confusing, imprecisely described processes that go under the name of measurement, requiring the discretion of the theoretical physicist for their proper interpretation, can all be unambiguously reduced to elementary 1-Qbit measurement gates, acting in conjunction with unitary gates, though often not without some quantum-computational programming ingenuity. I give a few examples of this below.

Figure 2.

The action of the 1-Qbit measurement gate shown in Figure 1 is actually a special case of its more general action, when applied to one of $N$ Qbits, initially in some $N$-Qbit state $|\Psi\rangle$. Figure 2 gives this more general definition of the 1-Qbit measurement gate. If $P_0$ and $P_1$ project a general $N$-Qbit state into its components along the states $|0\rangle$ and $|1\rangle$ in the subspace associated with the Qbit on which the measurement gate acts, then the general $N$-Qbit input state $|\Psi\rangle$ can be expanded as

$$|\Psi\rangle = (P_0 + P_1)|\Psi\rangle = a_0|0\rangle|\Phi_0\rangle + a_1|1\rangle|\Phi_1\rangle,$$

where $|\Phi_0\rangle$ and $|\Phi_1\rangle$ are normalized (but in general nonorthogonal) states of the $N - 1$ other Qbits and $|a_0|^2 + |a_1|^2 = 1$. If the measurement gate flashes $x$, then the $N$-Qbit state after the measurement is $|x\rangle|\Phi_x\rangle$ with probability $p_x = |a_x|^2$.
Born rule), is demonstrated in Figure 3. If the single Qbit subject to the measurement gate is unentangled with the other \( N - 1 \) Qbits, then \( |\Phi_0\rangle = |\Phi_1\rangle = |\Phi\rangle \). Consequently Figure 2 reduces to Figure 3. The outgoing state of the \( N - 1 \) unmeasured Qbits continues to be its incoming state \( |\Phi\rangle \), independent of the reading of the measurement gate, so the unmeasured Qbits remain unentangled with the measured Qbit. The state of the unmeasured Qbits is entirely unaffected by the measurement gate, while the final state of the measured Qbit is exactly as specified in Figure 1.

\[
|\Psi\rangle = \begin{cases} 
  a_0|0\rangle + a_1|1\rangle & \text{for } x \\quad \text{(upper half of Figure 4)} \\
  |\Phi\rangle & \text{for } P_x = |a_x|^2 
\end{cases}
\]

**Figure 3.**

You might think there was need for another kind of measurement hardware — a still more general measurement gate that acts on all \( N \) Qbits in the manner specified in the upper half of Figure 4. But in fact it can be shown from the generalized Born rule specified in Figure 3 that precisely this effect is achieved by applying \( N \) different 1-Qbit measurement gates to the \( N \) individual Qbits, illustrated for \( N = 3 \) in the lower half of Figure 4. The proof that this result always obtains, independent of the order in which the 1-Qbit gates are applied, can be extracted from the generalized Born rule in a straightforward (but irritatingly clumsy) way.

Finally, all these measurements are what one would call, more generally, “measurement in the computational (classical) basis”. A more general (von Neumann) measurement on \( N \) Qbits would associate states \( |\Phi_x\rangle \) with the out-going Qbits taken from some general complete set of orthonormal states. Such states can, however, be unitarily related to the computational-basis states: \( |\Phi_x\rangle = U|x\rangle \). Consequently the more general \( N \)-Qbit measurement gate \( MU \) can be constructed out of the \( N \)-Qbit computational-basis measurement gate \( M \) as pictured in Figure 5.

So all von Neumann measurements can be constructed out of unitary gates combined with one or more specimens of a single black-boxed piece of hardware: the 1-Qbit measurement gate. Such obscurity as there might be in the measurement process rests entirely in that single elementary gate, whose action is unambiguously specified in Figure 1. If there is a “measurement problem” in quantum computation, it is a “1-Qbit measurement gate problem”.

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Note that one can construct an elementary measurement apparatus — another on Bell’s list of bad words — out of a 1-Qbit measurement gate, an ancillary Qbit initially in the state $|0\rangle$, and a cNOT gate that couples the Qbit to be measured to the ancilla, as pictured in Figure 6. The measurement gate no longer acts on the measured Qbit, but only on the ancilla. Nevertheless, the resulting final state for the measured Qbit and the associated probabilities are exactly the same as they would be if the measurement gate acted on it directly, thanks to its coupling to the “apparatus” via the cNOT gate.

One can also remove the action of the measurement gate from the ancilla, by introducing a second ancilla, coupled to the first by a second cNOT gate, as pictured in Figure 7, and can continue in this way for many stages, as pictured in Figure 8. At some point the process has to stop with a measurement gate, if one is to extract any information from the computer, but one is perfectly free to introduce as many ancillary Qbits as one wishes, before reaching this stage. This corresponds to Bell’s “shifty split” between what
one chooses to call “the quantum system” and what one chooses to call “the measurement apparatus”, but there is no longer anything shifty about it. A measurement takes place if and only if there is a 1-Qbit measurement gate somewhere in the circuit. In the absence of a measurement gate, since cNOT is its own inverse, one can undo the whole process and get back to the initial state, as shown in Figure 9.

So it is through the readings of 1-Qbit measurement gates, and only through such readings, that one can extract information about a computation from a quantum computer. But this is not the only role measurement gates play in quantum computer science. Measurement gates have a second, equally crucial task to perform, which is insufficiently emphasized in most expositions of the subject.
\[ |\psi\rangle = a_0 |0\rangle + a_1 |1\rangle \]

Figure 7.

\[ p_x = |a_x|^2 \]

Figure 8.
Figure 9.

The action of all gates, whether they are unitary gates or measurement gates, is given by a rule that specifies the state associated with the output Qbits in terms of the state associated with the input Qbits. But how is a state to be associated with the Qbits at the very beginning of the computation, before any gates have acted? The simplest answer sufficient to make a coherent whole of the subject, and, I would maintain, the only answer that is fully formulated in mathematical terms, with nothing left to the discretion of the theoretical physicist, is this:

Take a Qbit off the shelf, and send it through a 1-Qbit measurement gate. If the display on the gate indicates 0, associate the state $|0\rangle$ with the Qbit emerging from the gate. If the display indicates 1, associate the state $|1\rangle$. Thus one can associate the state $|0\rangle$ with a Qbit by sending it through a 1-Qbit measurement gate and then sending it through a NOT gate $X$ if and only if the display on the measurement gate indicates 1. This is pictured in Figure 10.

Should the initial Qbit already have a state associated with it, or should it be a member of a larger group of Qbits which already have a state associated with them — i.e. should it be entangled with the other Qbits in the group — then Figure 10 is a consequence
of the generalized Born rule for 1-Qbit measurement gates pictured in Figure 2. If, on the other hand, one has no basis for associating any initial state either with the Qbit or with a larger group of Qbits to which it belongs — even a stochastic association weighted with probabilities (i.e. a mixed state) — then the rule illustrated in Figure 10 must be regarded as an additional property of 1-Qbit measurement gates, consistent with, but independent of the behavior specified in Figure 2:

The 1-Qbit measurement gate can thus be used to define what it means for a Qbit, about whose past history nothing whatever is known, to be in the state $|0\rangle$. One can view this role of the 1-Qbit measurement gate in “initializing a Qbit to the state $|0\rangle$” as a special kind of cNOT gate, in which the single Qbit functions as both control (the upper Qbit in Figure 1) and target (the lower Qbit in Figure 1). Such an auto-erotic cNOT gate, unlike a normal 2-Qbit cNOT gate, is not unitary, since the state associated with its output is $|0\rangle$, independent of whether the state associated with its input is $|0\rangle$ or $|1\rangle$.

Up until now I have used the term “state” without actually defining it in mathematical terms (with nothing left to the discretion of the theoretical physicist). But now we can give a precise definition:

The state $|x\rangle$ ($x = 0$ or 1) is associated with a Qbit if the Qbit, perhaps in association with other Qbits, has been subjected to a series of gates whose actions, together with the readings of any gates that are measurement gates, imply, by the rules of gate operation specified above, that the state associated with the Qbit is $|x\rangle$. 
A particularly simple example of those rules is that the state $|x\rangle$ is associated with a single Qbit emerging from a 1-Qbit measurement gate that reads $x$.

More generally, to associate the state $|\Psi\rangle$ with $N$ Qbits means nothing more nor less than that the Qbits, perhaps in association with other Qbits, have been subjected to a series of gates whose actions, together with the readings of any measurement gates, imply that the state associated with the Qbits is $|\Psi\rangle$.

The reader may have noted my practice of replacing the customary phrase “state of the Qbits” with “state associated with the Qbits”. The reason for this awkward expansion is that the simpler phrase, to the extent that it suggests that the state resides in or is an inherent property of the Qbits, is incorrect. Given full access to the Qbits, but no information about them, there is, famously, nothing you can do to them to inform you of their associated state. That state is nothing more than a compact summary of all features of their past history relevant to determining the probabilities of the readings of subsequent measurement gates, applied either immediately or after the action of additional unitary gates.

Unlike the misleading phrase “state of the Qbits”, the phrase “state associated with the Qbits” immediately raises the question “associated by whom?” Quantum computer science provides a straightforward answer to this question: the user. Here “user” is the standard computer-science user, most commonly encountered in the term “user-friendly”. It should be clear that “user” does not mean God (who already knows the factors of all integers and therefore does not need to use a quantum computer to help him out). Nor does it mean the mouse that Einstein worried about, since no mouse (even Einstein’s) has any interest in factoring large integers. “The user” is nothing more or less than the user.

Having made this important point, I will now use the more graceful phrase “state of the Qbits” with the understanding that “of” is simply a compact abbreviation of “associated with” or “assigned to”.

Figures 12 and 13 explore the question of whether it might be possible to eliminate this use of measurement gates in assigning the state $|0\rangle$ to an initially stateless Qbit. The Qbit in question is associated with the middle wire in Figure 12. Its initial statelessness is taken to arise from its entanglement with many other “external” Qbits, associated with the heavier bottom wire. The entanglement is characterized by a state $a_0|0\rangle|\Phi_0\rangle + a_1|1\rangle|\Phi_1\rangle$, where $|0\rangle$ and $|1\rangle$ are associated with the Qbit and $|\Phi_0\rangle$ and $|\Phi_1\rangle$ are associated with the Qbits with which it is entangled.

Rather than applying a measurement gate directly to the Qbit to associate with it the state $|0\rangle$ as in Figure 10, we introduce an ancilla (top wire in Figure 12) coupled to the Qbit with a cNOT gate, as in the measurement apparatus of Figure 6. (Throughout the rest of this discussion I refer to the Qbit on which the state is to be conferred as “the Qbit” and the ancillary Qbit, as “the ancilla.”) By applying or not applying $X$ to the Qbit, depending on whether or not the measurement gate applied to the ancilla indicates 1 or 0, we can associate with the Qbit the state $|0\rangle$ even though no measurement gate has
been applied to it directly, though, of course, a measurement gate (applied to the ancilla) has still played a crucial role in this state assignment.

We can automate this process, replacing the NOT gate that acts or doesn’t act on the Qbit, depending on the reading of the measurement gate, by a cNOT controlled by the ancilla and targeted on the Qbit. This is pictured in the lower part of Figure 12 and the upper part of Figure 13.

But now we can use the (easily verified) fact that the combined outcome of a cNOT gate followed by a measurement gate on the target Qbit is unaltered by changing the
order in which the two gates act. This gives us the middle part of Figure 13, in which the measurement gate only acts after the Qbit has been assigned the state $|0\rangle$. The measurement gate (and the NOT gate that follows it if the reading of the measurement
gate is 1) can therefore be dropped, without altering the state assignment to the Qbit, leaving us with the bottom part of Figure 13. So we have managed to associate the state $|0\rangle$ with the Qbit, without having to make any use of a measurement gate. The measurement gate has been eliminated from the act of state preparation.

But to do this we required an ancilla in the state $|0\rangle$. Indeed, all the circuit at the bottom does is swap the roles played by the ancilla and the Qbit. The Qbit acquires the state $|0\rangle$ initially possessed by the ancilla, and the ancilla ends up entangled with the external Qbits in exactly the same way that the Qbit originally was. If we wish to restore the ancilla to its initial state, so that we have actually produced an additional Qbit associated with the state $|0\rangle$ rather than merely swapped the state $|0\rangle$ from one Qbit to another, then we must retain the final operations on the ancilla pictured in the middle part of Figure 13.

The situation here is reminiscent of Charles Bennett’s solution [3] to the problem posed by Maxwell’s demon. In the bottom part of Figure 13 we do indeed manage to get the Qbit into the pure state $|0\rangle$, but at the price of removing from the demon (ancilla) its own pure state. To get the demon back to a condition in which it can continue its work, it is necessary to erase the entanglement it has acquired with the external Qbits (by measuring it), and we have therefore failed in our attempt to eliminate measurement gates from the act of state preparation.

It is much the same with error correction, which is, in a sense, nothing more than a refined form of state preparation. To illustrate this, and to elaborate on yet another way in which measurement is unproblematic (as well as crucial) in quantum computation, it is necessary to say a little about the measurement of degenerate observables.

We have already seen an example of measuring a degenerate observable in the generalized Born rule, illustrated in Figure 2, the observable being either of the projection operators $P_x$ on the computational basis states $|x\rangle$ of the measured Qbit, each of which is $2^{N-1}$-fold degenerate in the $2^N$-dimensional space of $N$ Qbits. Devising ways to measure less trivial degenerate observables with 1-Qbit measurement gates can require a certain amount of quantum programming ingenuity, but this is surely what Bell meant when he deplored leaving things to the discretion of the theoretical physicist.

Figure 14 shows, as an important example, the degenerate 2-Qbit parity-measurement gate $P$, which outputs, with the appropriate probabilities, the normalized projection of the input 2-Qbit state into the even and odd parity subspaces, spanned by $|00\rangle$ and $|11\rangle$ (even) or $|10\rangle$ and $|01\rangle$ (odd). Figure 15 shows how to realize this gate using a single 1-Qbit measurement gate and the aid of an ancilla, as usual initialized to the state $|0\rangle$, which is coupled by cNOT operations to each of the two Qbits whose parity is being measured. Since the combined effect of the two cNOT gates on the ancilla is to leave it in the state $|0\rangle$ if the two measured Qbits are in either of the even parity computational basis states, and flip it to the state $|1\rangle$ if they are in either of the odd states, it follows from linearity and the generalized Born rule that this circuit achieves the parity measurement of Figure
as summarized in Figure 16. (The two final cNOT gates act to restore the ancilla to its initial state.)

This is one of the simplest examples of measuring a degenerate observable using ancilllas and 1-Qbit measurement gates. A more subtle example is shown in Figure 17. The upper part shows a circuit that measures the total spin of two spin-$\frac{1}{2}$ Qbits (with the understanding that the 1-Qbit state $|0\rangle$ represents spin up, and $|1\rangle$, spin down). The central 3-Qbit gate is a doubly-controlled NOT (or Toffoli) gate, which, acting on the computational basis, flips the state of the target Qbit if and only if both control Qbits are in the state $|1\rangle$. In quantum-information-theoretic language, the post-measurement state associated with the two Qbits is the anti-symmetric Bell state $\frac{1}{\sqrt{2}}(|0\rangle|1\rangle - |1\rangle|0\rangle)$ (total spin 0) if the reading $s$ of the measurement gate is 0, and is the projection of its initial state into the 3-dimensional symmetric subspace (total spin 1), if the measurement gate indicates 1.

To see that the circuit behaves in this way, note first that if $A$ is any $N$-Qbit unitary gate, and an ancilla acts as a control Qbit for a controlled-A gate, then if the state of the ancilla is $H|1\rangle$ prior to the action of the controlled-A gate, where $H$ is the Hadamard gate
\[ a_{00}|00\rangle + a_{01}|01\rangle + a_{10}|10\rangle + a_{11}|11\rangle \]

\[
|0\rangle \left[ a_{00}|00\rangle + a_{11}|11\rangle \right] + |1\rangle \left[ a_{01}|01\rangle + a_{10}|10\rangle \right] = \\
\sqrt{p_0} \left[ |0\rangle \frac{a_{00}|00\rangle + a_{11}|11\rangle}{\sqrt{p_0}} \right] + \sqrt{p_1} \left[ |1\rangle \frac{a_{01}|01\rangle + a_{10}|10\rangle}{\sqrt{p_1}} \right]
\]

Figure 15.

of Figure 1, and a second Hadamard gate is applied to the ancilla after the controlled-A
gate acts, then the final state of the $N+1$ Qbits is easily confirmed to be the one shown in the lower part of Figure 17.

$$|1\rangle \rightarrow H \rightarrow \begin{array}{c} \text{spin } s \\ \{ \text{part of } |\Psi\rangle \} \end{array} \rightarrow H \rightarrow |s\rangle$$

Figure 17.

The operation between the two Hadamards in the upper part of Figure 17 is precisely of this form (for $N=2$) with $A$ being a set of three cNOT gates (with alternating control and target Qbits) acting on the lower two Qbits. Such a set of three alternating cNOT gates is easily confirmed to be the SWAP gate, which, acting on the computational basis state $|x\rangle|y\rangle$ gives $|y\rangle|x\rangle$. But when $A$ is a two-Qbit swap gate, then $\frac{1}{2}(1-A)$ projects on the (1-dimensional) antisymmetric subspace, while $\frac{1}{2}(1+A)$ projects on the (3-dimensional) symmetric subspace.
Therefore the generalized Born rule requires the final state associated with the two lower Qbits to be the symmetric or anti-symmetric component of the input state, depending on whether the measurement gate acting on the ancilla reads 0 or 1.

An amusing feature of this example (and many other quantum computational circuits) is that there is quite a different way to understand why it behaves as it does. In this case one can construct an alternative view of the circuit by noting that if any control Qbit in a multiply-controlled NOT gate is sandwiched between Hadamards, then, without changing the action of the circuit, the control Qbit can be changed to a target Qbit (without the two Hadamards) provided the former target Qbit is changed to a control Qbit sandwiched between two Hadamards. (This follows most simply from the fact that $X = HZH$, where $Z$ is the phase gate, $Z|x⟩ = (−1)^x|x⟩$, and the fact that the action of a multiply-controlled $Z$ gate is unaltered by interchanging the target Qbit with any of the control Qbits, since all it does is multiply the $(N + 1)$-Qbit computational basis state by $−1$ if and only if that states is a product of $N + 1$ states $|1⟩$.)

\[
\begin{align*}
|1⟩ & \rightarrow \circ \rightarrow \circ \rightarrow |s⟩ \\
|Ψ⟩ \{ & \text{Bell-basis to computational basis.} \\
H & \rightarrow \circ \rightarrow H \rightarrow \circ \\
& \text{Computational basis to Bell basis}
\end{align*}
\]

**Figure 18.**

Figure 18 shows the upper part of Figure 17, modified in precisely this way. But the action of this circuit has quite a different interpretation. The first two gates on the left act on the lower two Qbits to convert the Bell basis to the computational basis. In particular
the (singlet) Bell state \( \frac{1}{\sqrt{2}} (|0\rangle|1\rangle - |1\rangle|0\rangle) \) is taken into the computational basis state \( |1\rangle|1\rangle \), while the three symmetric (spin 1) Bell states are taken into the computational basis states \( |1\rangle|0\rangle, |0\rangle|1\rangle, \) and \( |0\rangle|0\rangle \). So if \( |\Psi\rangle \) is the singlet state, then the NOT acts on the ancilla and the measurement gate registers 0, but if \( |\Psi\rangle \) is any triplet (spin-1) state, the NOT does not act and the measurement gate registers 1.

I conclude with a glimpse into the role played by measurement in error correction, emphasizing the conceptual similarity to state preparation. Suppose we encode any 1-Qbit state \( \alpha|0\rangle + \beta|1\rangle \) into a 3-Qbit code state \( \alpha|0\rangle|0\rangle|0\rangle + \beta|1\rangle|1\rangle|1\rangle \) and suppose (artificially) that the only corruption that can befall the three Qbits is that at most one of them may be subject to a NOT gate (which is the general classical single-Cbit error, but not the general quantum single-Qbit error.) We require a circuit that tests for such an error and, if a bit-flip error is found, restores the original encoded state \( \alpha|0\rangle|0\rangle|0\rangle + \beta|1\rangle|1\rangle|1\rangle \), whatever the values of the amplitudes \( \alpha \) and \( \beta \).

Figure 19 shows the entire process of error production and correction. On the lower left we have the unencoded 1-Qbit state \( \alpha|0\rangle + \beta|1\rangle \) and two additional Qbits in the state \( |0\rangle \). The two initial cNOT gates on the lower left act to transform these three Qbits into the 3-Qbit code state \( \alpha|0\rangle|0\rangle|0\rangle + \beta|1\rangle|1\rangle|1\rangle \). The 3 ghost-like NOT gates in noisy-looking boxes indicate the possible errors. At most one of them is realized as an actual NOT gate \( X \). It is the job of the rest of the circuit to determine which, if any, of the Qbits suffered such an extraneous NOT, and to restore the initial 3-Qbit code state \( \alpha|0\rangle|0\rangle|0\rangle + \beta|1\rangle|1\rangle|1\rangle \) regardless of which (if any) of the Qbits was so corrupted.

The next two cNOTs, that target the upper ancilla, together with the measurement gate on the upper ancilla, act as the parity-measuring circuit of Figure 15, as do the final two cNOTs targeting the lower ancilla and the other measurement gate. Since the 3 Qbits in both basis states for the code-word subspace have the same parity, the measurement gate on the upper ancilla will indicate \( x = 1 \) if and only if either of the upper two code-word Qbits suffered a NOT error, while the measurement gate on the lower ancilla will indicate \( y = 1 \) if and only if either of the lower two-code-word Qbits suffered such an error.

Consequently if both measurements indicate 0, then there was no error, and no correction to be applied. If the upper measurement gate indicates 1 and the lower indicates 0, the upper code-word Qbit must have suffered the error, and applying a NOT gate to it will restore the original code-word state. If the upper measurement gate indicates 0 and the lower indicates 1, the lower code-word Qbit must have suffered the error, and the original code-word will be restored by applying the NOT gate to the lower code-word Qbit. And if both measurement gates indicate 1 it is the middle code-word Qbit that suffered an error, and to which the corrective NOT gate must be applied. This is precisely what the NOT gates on the right of Figure 19 accomplish.

How essential are the two measurement gates to this process? Can the corrective NOT gates be applied by judicious use of unitary controlled gates, without having to perform any measurements at all? Figure 20 shows some steps in this direction.
The upper part of Figure 20 reproduces the relevant part of Figure 19. The middle part replaces the action of the three corrective NOT gates, contingent on the three measurement outcomes, by a collection of unitary gates — two cNOT gates and a doubly-controlled triple-NOT gate (that can be built out of three Toffoli gates) — that are controlled by the post-measurement ancillary Qbits, and are easily confirmed to apply just the required

\[ a = x + xy \]
\[ b = xy \]
\[ c = y + xy \]
or

\( \alpha |000\rangle + \beta |111\rangle \)

or

\( x^a \)  \( (a = x + xy) \)

or

\( x^b \)  \( (b = xy) \)  \( \alpha |000\rangle + \beta |111\rangle \)

or

\( c = y + xy \)

Figure 20.

corrections.
Now, just as in Figure 13, we may let these controlled gates act prior to the measurement gates acting on their Qbits, without altering either the outcomes or the probabilities of those outcomes. This is pictured in the lower part of Figure 20.

But now the measurement gates act on the two ancillary Qbits only after the error has been fully diagnosed and corrected by unitary gates. If we drop them from the circuit, as pictured in Figure 21, the error has been corrected without using any measurement gates at all.

The only problem with this circuit is the one we faced in trying to produce a Qbit in the state $|0\rangle$, without using a measurement gate. There the measurement gate was needed to restore the ancilla to the state $|0\rangle$ so it could be used again. It is the same for error correction. Having removed the measurement gates from the ancillas, we no longer know whether their states are $|0\rangle|0\rangle$ (corresponding to no error) or $|1\rangle|0\rangle, |0\rangle|1\rangle$ or $|1\rangle|1\rangle$ (corresponding to an error in the top, bottom, or middle code-word Qbit). To get ancillary Qbits for the next round of error correction we must either initialize two fresh Qbits to the state $|0\rangle|0\rangle$ by means of measurement gates, or measure the two we already have,
to determine which if any must be flipped to restore their initial state. In either case, measurement gates must be employed.

So, in summary, measurement is essential at both ends of a quantum computation, nor is it problematic in quantum computer science. It’s problematic character is softened because the state assigned to the Qbits is changed discontinuously by all gates — not just by measurement gates — thereby eliminating the disconnect between “continuous Schrödinger evolution” and “discontinuous collapse”. To be sure, there remains the distinction between linear, invertible, unitary gates with output state fully determined by the input, and nonlinear, irreversible, measurement gates, with output state stochastically determined by the input. But the actions of both types of gates are fully defined, with nothing left to the discretion of the theoretical physicist.

Modulo unitary gates, all the diverse things that go under the name of “measurement” reduce to (multiple copies of) a single elementary 1-Qbit gate. If there remains anything problematic about measurement in quantum computer science, it lies entirely in the straightforward action of the simple circuit pictured at the bottom of Figure 1. And because all measurements must be constructed out of 1-qubit gates, the complex network of interactions (unitary gates — often cNOT gates or multiply controlled NOT gates) hidden inside many superficially elementary measurements is always made explicit. Bell’s question, — what exactly qualifies as a measurement — has a precise and unambiguous answer. A circuit may qualify as a measurement if and only if contains at least one 1-Qbit measurement gate.

In quantum computation, there is a system (one more of Bell’s bad words) — the computationally relevant Qbits — and there are things outside of it — the gates that act on those Qbits, the readings of the measurement gates, and the users who read those readings. There is no point to a computer unless users learn something from it. There must be coupling to something on the outside (users) if the computer is to perform its function.

Does this generalize beyond quantum computation? Unlike a computer, the universe is not made just for us. But theoretical physics is made just for us. There is no point to quantum mechanics unless we learn something from it. There must be coupling to something on the outside (us) if the quantum theory is to perform its function. This is what I take to be the meaning of the Bohr quotation at the head of this paper: “In our description of nature the purpose is not to disclose the real essence of the phenomena but only to track down, so far as it is possible, relations between the manifold aspects of our experience.” Our description of nature is theoretical physics — quantum mechanics, in particular. It is not concerned with “real essences” and whether or not that term has any meaning, it certainly does not within the setting of quantum mechanics. What does have meaning and what science strives to do is to impose an orderly structure on our perceptions.
Compare quantum computation: in the theory of quantum computation the purpose is not to disclose the real essence of the Qbits, but only to track down correlations — brought about by the action of intermediate gates — between our initial and final readings of the elementary 1-Qbit measurement gates.

The world is, of course, more complex than a quantum computation. A computation uses a finite set of Qbits. It has an unambiguous beginning and end. There is always a world external to the computation. If there were not an outside world, there would be no point in doing the computation because there would be nobody or nothing to take advantage of the output.

Nobody (well practically nobody) wants to view the entire universe as one colossal quantum computer, sufficient unto itself. We should adopt a similar modesty of scope in our view of quantum mechanics. Physics is a tool for relating some aspects of our experience to other aspects. Every application of physics can be cast in a form that begins and ends with statements about experience. This (and not the existence of parallel computations in parallel universes) is the great lesson that quantum computer science teaches us.

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