On the Lie complexity of Sturmian words

Alessandro De Luca

DIETI, Università degli Studi di Napoli Federico II, Italy

Gabriele Fici

Dipartimento di Matematica e Informatica, Università di Palermo, Italy

Abstract

Bell and Shallit recently introduced the Lie complexity of an infinite word $s$ as the function counting for each length the number of conjugacy classes of words whose elements are all factors of $s$. They proved, using algebraic techniques, that the Lie complexity is bounded above by the first difference of the factor complexity plus one; hence, it is uniformly bounded for words with linear factor complexity, and, in particular, it is at most 2 for Sturmian words, which are precisely the words with factor complexity $n + 1$ for every $n$. In this note, we provide an elementary combinatorial proof of the result of Bell and Shallit and give an exact formula for the Lie complexity of any Sturmian word.

Keywords: Sturmian word, Lie complexity.

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1. Introduction

The factor complexity $p_w$ of an infinite word $w$ is the integer function that counts, for every nonnegative integer $n$, the number of distinct factors of length $n$ occurring in $w$. This notion is widely used in the combinatorial investigation of infinite sequences. For example, it is used in the definition of topological entropy of a symbolic dynamical system.

A fundamental result of Morse and Hedlund [13] is that any aperiodic right-infinite word has factor complexity at least $n + 1$ for every $n$. Sturmian words are aperiodic words with minimal factor complexity, i.e., they have factor complexity equal to $n + 1$ for every $n$ (in particular they have two factors of length 1, i.e., they are binary words).

In the literature, other complexity functions have been introduced. To cite a few, abelian complexity [7], $k$-abelian complexity [10], arithmetic complexity [1, 5], maximal pattern complexity [9], cyclic complexity [4], binomial complexity [15], window complexity [6], periodicity complexity [12], etc.

*Corresponding author.

Email addresses: alessandro.deluca@unina.it (Alessandro De Luca), gabriele.fici@unipa.it (Gabriele Fici)

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Recently, Bell and Shallit introduced the notion of Lie complexity of an infinite word \( w \) as the integer function whose value at \( n \) is the number of conjugacy classes (under cyclic shift) of factors of length \( n \) of \( w \) with the property that every element in the conjugacy class occurs as a factor in \( w \). We call such a conjugacy class a Lie class of factors of \( w \).

Bell and Shallit proved the following result:

**Theorem 1** ([3]). Let \( \Sigma \) be a finite alphabet, let \( w \) be a right-infinite word over \( \Sigma \), and let \( L_w : \mathbb{N} \rightarrow \mathbb{N} \) be the Lie complexity function of \( w \). Then for each \( n \geq 1 \) we have

\[
L_w(n) \leq p_w(n) - p_w(n - 1) + 1.
\]

Hence, the Lie complexity is uniformly bounded for words with linear factor complexity, and, in particular, it is bounded by 2 for Sturmian words.

The proof of the previous theorem given in [3] is purely algebraic. In this note, we provide an elementary combinatorial proof of this result.

We then give an exact formula for the Lie complexity of any Sturmian word of slope \( \alpha \) in terms of the continued fraction expansion of \( \alpha \). For a general introduction to Sturmian words the reader is pointed to [11].

### 2. A combinatorial proof for the bound on the Lie complexity

For all \( n \geq 0 \), let \( \text{Fact}_w(n) \) denote the set of factors of length \( n \) of \( w \), so that \( p_w(n) = \# \text{Fact}_w(n) \). Recall that the Rauzy graph of order \( n \geq 1 \) for \( w \), denoted by \( \Gamma_w(n) \), is the directed graph with set of vertices \( \text{Fact}_w(n - 1) \) and set of edges \( \text{Fact}_w(n) \) such that an edge \( e \in \text{Fact}_w(n) \) starts at vertex \( v \) and ends at a vertex \( v' \) if and only if \( v \) is a prefix of \( e \) and \( v' \) is a suffix of \( e \).

Recall that in a directed graph, a (simple) cycle is a walk that starts and ends in the same vertex and no other vertex is repeated. For our purposes, we identify cycles having the same sets of vertices (and edges).

**Lemma 2.** Lie classes of factors of length \( n \) correspond exactly to cycles whose lengths divide \( n \) in \( \Gamma_w(n) \).

**Proof.** Suppose that all cyclic shifts of \( u = a_1 \cdots a_n \) are factors of \( w \). Then such shifts correspond to consecutive edges of a cycle in the Rauzy graph; if they are all distinct, i.e., \( u \) is primitive, then clearly the cycle has length \( n \). Otherwise we can write \( u = v^{n/d} \) for some \( v \), with \( d \) the smallest index such that \( u = a_{d+1} \cdots a_n a_1 \cdots a_d \).

Conversely, let \( u_1, \ldots, u_d \) be consecutive edges of a cycle, with \( d|n \), and set

\[
\begin{align*}
  u_1 &= a_1 a_2 \cdots a_{n-1} x_1, \\
  u_2 &= a_2 \cdots a_{n-1} x_1 x_2, \\
  & \vdots \\
  u_d &= a_d \cdots a_{n-1} x_1 \cdots x_d
\end{align*}
\]

for letters $a_1, \ldots, a_{n-1}$ and $x_1, \ldots, x_d$. Since the last edge $u_d$ returns to the starting vertex $a_1 \cdots a_{n-1}$, the word $a_1 \cdots a_{n-1}x_1 \cdots x_d$ has $a_1 \cdots a_{n-1}$ as a suffix as well as a prefix. This implies that all its factors $u_1, \ldots, u_d$ have $d$ as a period, so that they are all the cyclic shifts of $u_1$.

In view of the previous lemma, we say that a cycle in the Rauzy graph $\Gamma_w(n)$ is a Lie cycle if its length divides $n$. Thus, $L_w(n)$ is the number of Lie cycles in $\Gamma_w(n)$, whereas $p_w(n)$ and $p_w(n-1)$ are the numbers of edges and vertices, respectively.

For a vertex $v$, we let $\text{odeg}(v)$ denote the out-degree of $v$, i.e., the number of distinct edges leaving $v$.

Proof of Theorem 1. We first observe that two Lie cycles may share one or more vertices but cannot share edges, since conjugacy classes are disjoint. As a consequence, if a vertex belongs to $k$ different Lie cycles, its out-degree is at least $k$.

We show that in $\Gamma_w(n)$ there exists a set $\mathcal{L}$ of $L_w(n) - 1$ edges such that every vertex of $\Gamma_w(n)$ has an outgoing edge not belonging to $\mathcal{L}$; this proves that the number of edges minus the number of vertices is at least $L_w(n) - 1$, whence the claimed inequality $L_w(n) \leq p_w(n) - p_w(n-1) + 1$.

Consider a walk on $\Gamma_w(n)$ visiting at least one edge for each Lie cycle ($w$ itself provides an example of such a walk). With the possible exception of the last one visited, every Lie cycle must contain a vertex with out-degree at least 2. Suppose $v$ is such a vertex, and let $k \geq 1$ be the number of Lie cycles containing $v$, so that $\text{odeg}(v) \geq k$. Then, since the walk visits all Lie cycles in $\Gamma_w(n)$, at least one of the following cases occurs:

1. one of the $k$ cycles is the last one visited by the walk;
2. $\text{odeg}(v) \geq k + 1$;
3. at least one of the $k$ cycles contains a vertex $v' \neq v$ with $\text{odeg}(v') \geq 2$.

Therefore, we can define $\mathcal{L}$ as follows: for each of the first $L_w(n) - 1$ Lie cycles, we choose an edge belonging to the same Lie cycle and leaving from a vertex with out-degree at least 2, with the requirement that each of these vertices has at least one outgoing edge which is not chosen. This choice for $\mathcal{L}$ ensures that each vertex in $\Gamma_w(n)$ has at least one outgoing edge not belonging to $\mathcal{L}$, as required.

3. A formula for the Lie complexity of Sturmian words

A Sturmian word $s = s_{\alpha,\rho}$ over $\Sigma = \{0, 1\}$ can be defined by taking an irrational number $0 < \alpha < 1$ (called slope) and a real number $\rho$ (called intercept) and defining for each $n \geq 0$

$$s_{\alpha,\rho}(n) = \lfloor \alpha(n+1) + \rho \rfloor - \lfloor \alpha n + \rho \rfloor$$

As is well known, any two Sturmian words $s = s_{\alpha,\rho}$ and $s' = s'_{\alpha',\rho'}$ with the same slope have the same factors. Therefore, one often considers the characteristic Sturmian word of slope $\alpha$, which is the word $s_{\alpha,\alpha}$.
Let \([0; d_1 + 1, d_2, \ldots, d_n, \ldots]\) be the continued fraction expansion of \(\alpha\). We will assume that 11 is not a factor of \(s_{\alpha, \alpha}\), which corresponds to assuming \(d_1 > 0\), i.e., \(\alpha < 1/2\). The other case, i.e., when 11 is a factor of \(s_{\alpha, \alpha}\) for \(\alpha = [0; 1, d_2, d_3, \ldots]\), can be reduced to considering the characteristic Sturmian word obtained by exchanging the two letters, which has slope \(\alpha' = [0; d_2 + 1, d_3, \ldots]\).

The characteristic Sturmian word \(s = s_{\alpha, \alpha}\) is the limit of the sequence of finite words \(s_{-1} = 1, s_0 = 0\) and \(s_n = s_{n-1}^{d_n} s_{n-2}\) for \(n > 0\). The words \(s_k, k \geq 0\), are called standard prefixes of \(s\).

For each \(k \geq 0\), the length of \(s_k\) is equal to \(q_k\), the denominator of the \(k\)-th convergent \(p_k/q_k = [0; d_1 + 1, d_2, \ldots, d_k]\) (we assume \(q_0 = 1\)). We will also need, when \(d_k > 1\), the denominators \(q_{k, \ell}\) of the \(k\)-th semiconvergents \(p_{k, \ell}/q_{k, \ell} = [0; d_1 + 1, d_2, \ldots, d_{k-1}, \ell], 1 \leq \ell < d_k\). The words \(s_{k, \ell} = s_{k-1 \ell} s_{k-2} \) of length \(q_{k, \ell}\) are sometimes called semistandard prefixes of \(s\).

Let \(S\) denote the set of standard or semistandard prefixes of \(s\). For every word \(v \in S\) of length at least 2, one has \(v = uab\), where \(ab \in \{01, 10\}\) and the word \(u\), called a central prefix, is a bispecial factor of \(s\). Recall that a factor \(u\) of \(s\) is left (resp. right) special if both \(0u, lu\) (resp. both \(u0, ul\)) are factors of \(s\) and bispecial if it is both left special and right special. Notice that since a Sturmian word has \(n + 1\) factors of length \(n\), it must have exactly one left (resp. right) special factor of each length \(n\), and this must therefore be a prefix (resp. suffix) of a bispecial factor.

The following result follows from [4, Lemma 9].

**Lemma 3.** Let \(s\) be a Sturmian word and \(w\) a primitive factor of \(s\) of length at least 2. Then all conjugates of \(w\) are factors of \(s\) if and only if \(w\) is a conjugate of an element of \(S\).

The best known example of a Sturmian word is the Fibonacci word \(f = 0100101001001 \cdots\), which can be defined as the fixed point of the morphism sending 0 to 01 and 1 to 0. The Fibonacci word is intimately related to the well-known sequence of Fibonacci numbers: \(F_1 = 1, F_2 = 1, F_n = F_{n-1} + F_{n-2}\) for \(n \geq 2\). The precise relation is the following: \(f\) is the characteristic Sturmian word \(s_{1/\phi^2, 1/\phi^2}\), where \(\phi = (1 + \sqrt{5})/2\) is the golden ratio. Since \(1/\phi^2 = [0; 2, 1]\), we have that for the Fibonacci word \(d_n = 1\) for every \(n\) and the sequence \(q_n = F_{n+2}\) is the sequence of denominators of the convergents of \(1/\phi^2\). The standard prefixes of \(f\) (of length \(F_n\)) are the Fibonacci finite words 1, 01, 010, 0101, etc.

In Example 7.4 of [3], the authors looked at the Lie complexity \(L_f\) of the Fibonacci word \(f\) and showed that

\[
L_f(n) = \begin{cases} 1, & \text{if } n = 0 \text{ or } n = F_k \text{ for } k \geq 4 \text{ or } n = F_k + F_{k-3} \text{ for } k \geq 4; \\ 2, & \text{if } n = 1, 2; \\ 0, & \text{otherwise}. \end{cases}
\]

Notice that \(F_k + F_{k-3} = F_{k-1} + F_{k-2} + F_{k-3} = 2F_{k-1}\).

The main result of this section is the following:
Theorem 4. The Lie complexity of any Sturmian word $s$ of slope $\alpha < 1/2$ is:

$$L_s(n) = \begin{cases} 
1, & \text{if } n = 0 \text{ or } n = q_{k,\ell} \text{ for } k \geq 2 \text{ or } n = mq_k \text{ for } 1 \leq m \leq d_{k+1} + 1 \text{ and } k \geq 1; \\
2, & \text{if } n = 1, 2, \ldots, q_1; \\
0, & \text{otherwise}.
\end{cases}$$

Lemma 5. Let $s$ be a Sturmian word and $w$ a factor of $s$ of length at least 2. If all conjugates of $w$ are factors of $s$, then $w$ is a power of a conjugate of an element of $S$.

Proof. If $w = v^m$, $v$ primitive, and all conjugates of $w$ are factors of $s$, then in particular all conjugates of $v$ are factors of $s$, hence by Lemma 3, $v$ is a conjugate of an element of $S$. \qed

Example 6. The converse is not true. Consider the Fibonacci word $f = 010010010100101001 \cdots$. The factor $w = (010)^3$ is a power of the standard prefix 010, yet no other conjugate of $w$ is a factor of $f$.

The following result is due to Damanik and Lenz [8, Thm. 4] (see also [14]). Recall that the index of a factor $v$ of $s$ is the largest integer $n$ such that $v^n$ is a factor of $s$.

Theorem 7 ([8]). Let $s$ be a Sturmian word.

- All conjugates of the standard prefix $s_1$ have index $d_2 + 1$;
- For every $k \geq 2$, the set of indexes of all conjugates of the standard prefix $s_k$ is $\{d_{k+1} + 1, d_{k+1} + 2\}$;
- For every $k \geq 2$, the set of indexes of all conjugates of a semistandard prefix $s_{k,\ell}$ is $\{1, 2\}$.

Corollary 8. Let $s$ be a Sturmian word.

- For every $k \geq 1$ and $1 \leq m \leq d_{k+1} + 1$, all conjugates of $s_k^m$ are factors of $s$, but not all conjugates of $s_k^{d_{k+1}+2}$ are factors of $s$;
- For every $k \geq 2$, all conjugates of $s_{k,\ell}$ are factors of $s$, but not all conjugates of $s_{k,\ell}^2$ are factors of $s$.

We are now able to give the proof of Theorem 4.

Proof. The assertion is trivially verified for $n = 0$, as well as for $1 \leq n \leq q_1$ since the $n + 1$ factors of $s$ of length $n$ are 0$^n$ and the $n$ conjugates of 0$^{n-1}$.

Let then $n > q_1$, and suppose $L_s(n) > 0$, so that there exists a factor $w$ of length $n$ such that all conjugates of $w$ are factors of $s$. By Lemma 3 there exists $v \in S$ and $m \geq 1$ such that all conjugates of $v^m$ are factors of $s$. Since $n > q_1$, either $v = s_{k-1}$ or $v = s_{k,\ell}$ for some $k \geq 2$. By Corollary 8 $n = q_{k,\ell}$ for $k \geq 2$ or $n = mq_k$ for $1 \leq m \leq d_{k+1} + 1$ and $k \geq 1$.\qed
To conclude the proof, we must show that $L_s(n) \leq 1$ for $n > q_1$, i.e., that the prefix $v$ is uniquely determined by $n$. For $k \geq 1$ and $1 \leq \ell < d_{k+1}$, by definition one has the following:

$$q_{k+1} = d_{k+1}q_k + q_{k-1}, \quad q_{k+1,\ell} = \ell q_k + q_{k-1}. \quad (1)$$

In particular, the sequence $(q_k)$ is strictly increasing; let then $k \geq 2$ be such that $q_k \leq n < q_{k+1}$, where $L_s(n) > 0$. By the above argument, the possible values for $n$ are

1. $mq_k$, for $1 \leq m \leq d_{k+1}$,
2. $(d_k + 1)q_{k-1} = q_k + q_{k-1} - q_{k-2}$,
3. $q_{k+1,\ell}$, for $1 \leq \ell < d_{k+1}$.

In view of (1), these are all distinct, so that the corresponding value for $|v|$ (respectively $q_k$, $q_{k-1}$, and $q_{k+1,\ell}$) is well defined and uniquely determined.

\[ \square \]

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