DISCRETIZATION STRATEGIES FOR COMPUTING CONLEY INDICES AND MORSE DECOMPOSITIONS OF FLOWS

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ABSTRACT. Conley indices and Morse decompositions of flows can be found by using algorithms which rigorously analyze discrete dynamical systems. This usually involves integrating a time discretization of the flow using interval arithmetic. We compare the old idea of fixing a time step as a parameters to a time step continuously varying in phase space. We present an example where this second strategy necessarily yields better numerical outputs and prove that our outputs yield a valid Morse decomposition of the given flow.

1. INTRODUCTION

While the numerical approximation of ordinary differential equations has a long history, the systematic study of how to compute time invariant structures began in the 1980s. Rigorous computations of these structures is an even more recent phenomenon (see [1, 8] and references therein). These latter efforts can be roughly divided into two approaches: direct computation of invariant sets, e.g., periodic orbits, heteroclinic and homoclinic orbits, invariant manifolds, and a more indirect approach based on identification of isolating neighborhoods. The advantages of the direct approach are clear: it leads to efficient algorithms and provides precise numerical bounds on solutions. The disadvantage is that, in general, it does not provide information about the global structure of the dynamics and the results are sensitive to changes in parameters. As indicated below, the indirect approach overcomes these disadvantages, however, the construction of appropriate methods of approximation has proven to be a significant challenge.

For the sake of simplicity consider a smooth ordinary differential equation

\[ \dot{x} = v(x), \quad x \in \mathbb{R}^d \]

that generates a flow \( \varphi : \mathbb{R} \times \mathbb{R}^d \rightarrow \mathbb{R}^d \). The focus of this paper is on obtaining a finite representation of \( \varphi \) from which information concerning the structure of the dynamics can be obtained. A theoretical resolution to this challenge is as follows.

Recall that given \( N \subset \mathbb{R}^d \) the maximal invariant set in \( N \) under the flow \( \varphi \) is

\[ \text{Inv}(N, \varphi) := \{ x \in N \mid \varphi(R, x) \subset N \}. \]

A compact set \( N \subset \mathbb{R}^d \) is an isolating neighborhood under \( \varphi \) if Inv\( (N, \varphi) \subset \text{int}N \). Then Inv\( (N, \varphi) \) is called an isolated invariant set. The Conley index (see Section [3,2]) is an algebraic topological invariant that provides information about the existence and structure of isolated invariant sets.

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To perform computations, we make use of a discretization in time. Choose \( h > 0 \) and define the diffeomorphism \( \phi_h : \mathbb{R}^d \to \mathbb{R}^d \) by
\[
\phi_h(x) := \varphi(h, x).
\]
In analogy with the setting of flows, the maximal invariant set in \( \mathbb{R}^d \) under \( \phi_h \) is defined by
\[
\text{Inv}(N, \phi_h) := \{ x \in N \mid \phi_h^n(x) \in N, \forall n \in \mathbb{Z} \}
\]
from which one can define isolating neighborhoods and isolated invariant sets.

A fundamental result [14, Theorem 1] guarantees that \( S \subset \mathbb{R}^d \) is an isolated invariant set for \( \varphi \) if and only if \( S \) is an isolated invariant set for \( \phi_h \), independent of the choice of \( h > 0 \). Furthermore, according to [14, Theorem 2] the Conley index of \( S \) computed using \( \phi_h \) determines the Conley index for \( S \) under \( \varphi \). Thus from a mathematical perspective no information is lost by studying the map \( \phi_h \) as opposed to the flow \( \varphi \).

To obtain a representation of \( \phi_h \) with which we can compute, we discretize the phase space. For the sake of simplicity of presentation, we assume that we are interested in understanding the global dynamics of \( \varphi \) or equivalently \( \phi_h \) restricted to a rectangular isolating neighborhood \( X \subset \mathbb{R}^d \). This allows us to discretize \( X \) using a cubical grid \( \mathcal{X} \) (see Section 2) with diameter less than some given \( \epsilon > 0 \). The dynamics of \( \phi_h \) can be encoded using a multivalued map \( \mathcal{F}_\rho : \mathcal{X} \rightarrow \mathcal{X} \) defined by
\[
\mathcal{F}_\rho(\xi) := \{ \xi' \in \mathcal{X} \mid B_\rho(\phi_h(\xi)) \cap \xi' \neq \emptyset \}
\]
where \( B_\rho(\phi_h(\xi)) := \{ x \in X \mid \inf_{y \in \phi_h(\xi)} \{ \| x - y \| \} < \rho \} \).

Making use of the fact that \( \mathcal{F}_\rho \) can be interpreted as a directed graph, a wide variety of efficient algorithms have been developed for finding isolating neighborhoods and Morse decompositions for \( \phi_h \) (see [2, 11]) and for computing the associated Conley indices [10, 15, 21]. Furthermore, by choosing finer grids, i.e., letting \( \epsilon \to 0 \) and making better approximations, i.e., letting \( \rho \to 0 \), one can recover any isolated invariant set [11]. More generally, one can find all attractor-repeller pairs or equivalently all Morse decompositions and in the limit recover Conley’s fundamental decomposition theorem [3, 11]. Thus, from a theoretical perspective this approach allows us to recover the local and global dynamics associated with isolated invariant sets generated by ordinary differential equations.

However, in concrete applications there are significant technical obstructions. Observe that there are three explicit parameters in the above mentioned approach: the time step \( h \), the grid size \( \epsilon \), and the accuracy of the approximation of the dynamics \( \rho \). Furthermore, the dynamics that can be extracted is quite sensitive to the particular choices of these parameters. For example, for a fixed \( \epsilon \), if \( h \) is too small then \( \xi \in \mathcal{F}_\rho(\xi) \) for every \( \xi \in \mathcal{X} \) and every \( \rho > 0 \). In this case no isolating neighborhood can be extracted and hence there is no interesting dynamics that can be resolved. Notice that the number of grid elements increases rapidly as an inverse function of \( \epsilon \) and thus it becomes computationally prohibitive to choose \( \epsilon \) too small. Observe that if \( v(x) \neq 0 \) for all \( x \in \xi \), then it is reasonable to assume that for a sufficiently large \( h \), \( \xi \notin \mathcal{F}_\rho(\xi) \), at which point one may hope to be able to start extracting interesting dynamical features. Clearly, the size of the parameter \( h \) depends on \( \|v\| \), the magnitude of the vector field. The naive conclusion is that \( h \) should be chosen to be large. However, this leads to other difficulties. We need to evaluate \( \phi_h \) which in practice requires us to integrate the differential equation (1) over time \( h \). Simple Gronwall inequalities imply that, in general, the numerical bounds on errors associated with this integration will grow exponentially in \( h \). Suggesting that large \( h \) leads to exponentially large
To address these issues, in this paper we consider a more general approach to encoding the dynamics. Let \( \tau : X \rightarrow (0, \infty) \) be a continuous function and consider the map \( \phi_\tau : X \rightarrow \mathbb{R}^d \) defined by

\[
\phi_\tau(x) := \phi(\tau(x), x).
\]

In principle this allows us to vary the integration time over the phase space, which allows us to compensate for the varying magnitude of the vector field and sensitivity of propagation of numerical errors. A brief outline of the paper is as follows.

Section 2 presents numerical examples that contrast results using fixed time steps \( \phi_h \) and spatially varying time steps \( \phi_\tau \). The goal is not to provide new results concerning the dynamics of a particular ordinary differential equation, but rather to demonstrate the potential usefulness of this approach.

We use the results of Section 2 as motivation for the central mathematical results of this paper. As indicated above, \([14, \text{Theorem 1}]\) guarantees that on the level of isolated invariant sets the dynamics of \( \phi_h \) captures the dynamics of \( \phi_\tau \). Since this is not always true for a \( \phi_\tau \) with varying time steps we prove Lemma 3.18, which provides a criterion under which we obtain the desired isolation.

In Section 3.2, we show that computing the Conley index of an isolated invariant set \( S \) for \( \phi_\tau \) and \( \phi \) provides the Conley index for \( S \) under the flow \( \phi \). This generalizes \([14, \text{Theorem 2}]\). In Section 3.3, we prove that a Morse decomposition for \( \phi_\tau \) is a Morse decomposition for the flow \( \phi \) if all the Morse sets are invariant under \( \phi \), thereby showing that our algorithm does compute what we are interested in.

Varying the time step in space was first proposed in \([7]\), but with less attention for the mathematical background, which we present here.

This article uses the following notations:

\[ \mathbb{R}_{\geq 0} := \{ x \in \mathbb{R} \mid x \geq 0 \}, \mathbb{R}_+ := \{ x \in \mathbb{R} \mid x > 0 \}. \]

We recall definitions where they are needed for our statements. Their proofs require some background in dynamical systems, which can be found in \([22]\) and other textbooks about the field.

### 2. A NUMERICAL EXAMPLE

#### 2.1. Combinatorial representations.

Given \( d > 0 \) and a tuple \( s \in \mathbb{R}^d_+ \) representing a box size, the space \( \mathbb{R}^d \) is the union of the following cubes (more precisely, products of intervals):

\[
\mathcal{Y} = \left\{ \prod_{i=1}^d [m_i s_i, (m_i + 1)s_i] \mid m \in \mathbb{Z}^d \right\}.
\]

Let \( \mathcal{X} \subset \mathcal{Y} \) be a finite subset. For a set \( \mathcal{A} \subset \mathcal{X} \), its realization is \( |\mathcal{A}| := \bigcup \{ \xi \mid \xi \in \mathcal{A} \} \subset \mathbb{R}^d \). Our goal is to compute a meaningful Morse decomposition within a compact set \( X := |\mathcal{X}| \).

A map \( \mathcal{F} \) from \( \mathcal{X} \) to its power set is written as \( \mathcal{F} : \mathcal{X} \Rightarrow \mathcal{X} \). We consider it as a multivalued map on \( \mathcal{X} \) or as a directed graph which has an edge from \( \xi \in \mathcal{X} \) to \( \xi' \in \mathcal{X} \) iff \( \xi' \in \mathcal{F}(\xi) \). The image of \( \mathcal{A} \subset \mathcal{X} \) is \( \mathcal{F}(\mathcal{A}) := \bigcup_{\xi \in \mathcal{A}} \mathcal{F}(\xi) \subset \mathcal{X} \). For a subset \( Z \subset X \), we let int\(X \) (\( Z \)) be the interior of \( Z \) with respect to the subspace topology of \( X \subset \mathbb{R}^d \).

**Definition 2.1.** A map \( \mathcal{F} : \mathcal{X} \Rightarrow \mathcal{X} \) is a combinatorial enclosure of \( f : \mathbb{R}^d \rightarrow \mathbb{R}^d \) if for every \( \xi \in \mathcal{X} : f(\xi) \cap X \subset \text{int}_X |\mathcal{F}(\xi)| \).
The equation has a fixed point \( (2) \) analyze with a fixed time step, 2.2. (ii) \( \text{Inv}(N,F) := \{ \xi \in N \mid \text{there is a full solution through } \xi \in N \} \).

We apply the algorithms and software described in [2] and [4] for finding Morse decompositions (Definition 3.17) as follows. The software constructs a combinatorial enclosure \( F : \mathcal{X} \xrightarrow{\text{in}} \mathcal{X} \) of a discrete dynamical system \( f : \mathbb{R}^d \rightarrow \mathbb{R}^d \). Then it constructs the strongly connected path components \( \{ \mathcal{N}(p) \mid p \in P \} \) of the directed graph \( F \) and builds the directed graph \( \text{MG}(F) \) with vertices \( P \) and a directed edge from \( p \) to \( q \) if there are \( G \in \mathcal{N}(p) \), \( H \in \mathcal{N}(q) \) and a path from \( G \) to \( H \) in \( F \). This describes the dynamics of the discrete dynamical system \( f \) because of the following Theorem 4.1 from [11].

**Theorem 2.3.** Let \( X = |\mathcal{X}| \) be an isolating neighborhood for \( f \) and \( S = \text{Inv}(X,f) \). Then:

(i) Each set \( |\mathcal{N}(p)| \) is an isolating neighborhood for \( f \).

(ii) \( \text{MG}(F) \) is a Morse graph for \( f \) in the invariant set \( S \) with Morse sets \( \text{Inv}(|\mathcal{N}(p)|,f), p \in P \).

In the rest of this section, we give an example flow \( \varphi \) and show that this algorithm yields a finer output when using \( f(x) = \varphi(\tau(x),x) \) than when using \( f(x) = \varphi(h,x) \). The justification that our outputs are indeed Morse decompositions for \( \varphi \) is formulated in Theorem 3.20. Every norm \( ||.|| \) in this article is Euclidean, i.e., \( ||x||^2 = \sum_{i=1}^{d} x_i^2 \) for \( x \in \mathbb{R}^d \).

2.2. **Fixed time step.** The following ordinary differential equation is particularly challenging to analyze with a fixed time step \( h \).

\[
\begin{align*}
\dot{x}_1 &= v_1(x) = -x_2 + x_1(2x_1^2 + x_2^2 - \mu)(x_1^2 + x_2^2 - 1) \\
\dot{x}_2 &= v_2(x) = x_1 + x_2(2x_1^2 + x_2^2 - \mu)(x_1^2 + x_2^2 - 1)
\end{align*}
\]

(2)

The equation has a fixed point \((0,0)\) and limit cycles with radius 1 and \( \sqrt{\mu} \) around the fixed point. This can be seen by its representation in polar coordinates:

\[
\begin{align*}
\dot{r} &= r(r^2 - \mu)(r^2 - 1) \\
\dot{\theta} &= 1
\end{align*}
\]

(3)

The norm of the vector field \( v \) increases quickly away from the origin because

\[
||v(x)|| = \sqrt{r^2(r^2 - \mu)^2(r^2 - 1)^2 + r^2} \approx r^5 \text{ for large } ||x||.
\]

Far away from the origin, the solutions behave like the solutions of \( \dot{r}(t) = r(t)^5 \). This equation is solved by

\[
\begin{align*}
r(t) &= \frac{1}{\sqrt{2} \sqrt[4]{(\sqrt{2} \cdot r_0)^{-4} - t}},
\end{align*}
\]

where \( r_0 = r(0) \). Note that the function \( r \) is only defined for \( t < (\sqrt{2} \cdot r_0)^{-4} = (x_1^2 + x_2^2)^{-2}/4 \). The trajectories reach infinity after finite time. For a point \( x \) with \( ||x|| > \sqrt{\mu} \), the solution for the system (2) is defined on a maximal open interval with right bound \( T_+(x) < \infty \). For \( h > T_+(x) \), the value \( \varphi_h(x) \) is undefined. Hence, also our integration algorithm fails when the input is a box
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Figure 1. Visualization of the flow from Equation (2) for $\mu = 2$ in $X = [-3,3] \times [-3,3]$ and outputs when using the fixed time step $h = 0.0006$ as described in Section 2.2. There are 57673 spurious combinatorial Morse sets. Most of them consist of just one box. Calculating the Morse graph was impossible because of memory problems.

containing $x$ and we ask it to integrate until time $h > T_+(x)$. The fixed time step strategy can only be used with a parameter $h < \min \{ T_+(x) \mid x \in X \}$. Therefore, when $X = [-3,3]^2$, we have to choose

$$h < T_+(3,3) \approx (3^2 + 3^2)^{-2}/4 = 1/1296 \approx 0.0007716.$$  

If $h$ is chosen larger, each box $\xi \in X$ near the boundary of $X$ is assigned an arrow in $F_h$ to all of the other boxes because the algorithm fails to find an enclosure of $\phi_h(\xi)$ (which does not even exist). This would lead to a very large invariant part $\text{Inv}(X, F_h)$ and should therefore be avoided. But when choosing $h$ small enough, the directed graph $F_h$ has a lot of small strongly connected components whose corresponding invariant set of $\phi$ is empty (so-called spurious Morse sets). They have to occur when the time step $h$ is so small that $\phi(h, \xi) \cap \xi \neq \emptyset$. But this happens easily near the origin where $\|v(x)\|$ is small.

The outputs for $\mu = 2$ are shown in Figure 1(b). The algorithm as described in [4] subdivided each dimension of $|X|$ into $2^9$ intervals of equal length. Each colored region is a combinatorial Morse set $N(p), p \in \mathcal{P}$. The index set $\mathcal{P}$ had 57675 elements. All but two of the combinatorial Morse sets found are empty.

It took 27 seconds on a laptop with an Intel i5 CPU to find the combinatorial Morse sets, but the Morse graph could not be computed because $\mathcal{P}$ was too large for the memory. Also subdividing each dimension into $2^{12}$ intervals did give similarly bad outputs (after 2396 seconds).

2.3. Variable time step. We use the following heuristic for the time step function $\tau$. For each subdivision level, the Euclidean norm $\|s\|$ is the diagonal of each of the congruent boxes $\xi \in \mathcal{X}$. Using parameters $D > 1$ and $\delta > 0$, we define the continuous function

$$\tau: X \to \mathbb{R}_+, \quad x \mapsto \frac{D\|s\|}{\|v(x)\| + \delta}.$$
The idea is that the distance between $x$ and $\varphi(\tau(x), x)$ should be around $D$ box diagonals – using a first-order approximation $\varphi(\tau(x), x) \approx x + \tau(x)v(x)$. The number $\delta$ ensures that $\tau$ is also defined when $v(x) = 0$. Since $\tau$ is usually not constant within a box $\xi$, the value $\tau(\xi)$ is an interval. We can find an enclosure of this interval by considering $\tau$ as a function on intervals and replacing $x$ by the box $\xi$ when calculating in interval arithmetic. The software library CAPD [5] is used to construct a combinatorial enclosure $F_\tau: \mathcal{X} \Rightarrow \mathcal{X}$ of $\varphi_{\tau}$. We use $D = 4$ and $\delta = 0.1$ in the following numerical example.

The varying time step strategy proposed above yields the finest Morse decomposition of the flow in $S = \text{Inv}([-3, 3]^2, \varphi) = \{x \mid \|x\| \leq \sqrt{11}\}$ (more precisely, the finest one which does not contain empty invariant sets). The output is shown in Figure 2.

Finding the Morse decomposition and the Morse graph took 29 seconds on the same hardware as in the first example, but with each dimension subdivided into only $2^8$ intervals. The algorithm additionally checks the correctness of the computed Morse decomposition. This verification and its output in Figure 2(c) are described in Remark 3.21.

3. Theoretical justification

For the remainder of the paper we let $Y$ be a locally compact separable metric space. Let $\varphi: \mathbb{R} \times Y \to Y$ be a flow. We show when certain invariants for $\varphi_{\tau}$ also yield the corresponding invariants for $\varphi$.

3.1. Isolating neighborhoods.

**Definition 3.1.** Let $f: Y \to Y$ be a continuous map.

(i) A solution through a point $x \in Y$ is a map $\gamma: \mathbb{Z} \to Y$ such that $\gamma(0) = x$ and $\gamma(n + 1) = f(\gamma(n))$ for all $n \in \mathbb{Z}$.

(ii) For $N \subseteq Y$ let $\text{Inv}(N, f) := \{x \in N \mid$ there is a solution $\gamma$ through $x$ such that $\gamma(\mathbb{Z}) \subseteq N\}$.

(iii) $N$ is an isolating neighborhood of the isolated invariant set $S$ if $S = \text{Inv}(N, f) \subseteq \text{int}N$.

Even though there is an inverse $\varphi_{-h}$ of $\varphi_h$ for any flow $\varphi$, there need not be an inverse for $\varphi_{\tau}$. 

**Figure 2.** Output for the same system as in Figure 1, but using time step function $\tau$ from Equation 4 as proposed in Section 2.3.
Lemma 3.2. Let $x \in Y$ and suppose that $\tau(\varphi([0,x])) \subset \mathbb{R}^+$ is bounded. Then there is a solution $\gamma: \mathbb{Z} \to Y$ of the discrete system $\varphi: Y \to Y$ such that $\gamma(0) = x$ and $\gamma(\mathbb{Z}) \subset \varphi([0,x])$.

Proof. Define $\gamma(n) = \varphi^n(x)$ for $n \geq 0$. Let $n < 0$ and assume that $\gamma(n+1)$ is already constructed. Then there is an $s \leq 0$ such that $\gamma(n+1) = \varphi(s,x)$. Define the function

$$g: \mathbb{R} \to \mathbb{R}, \quad t \mapsto \varphi(\varphi(t,x)) - s + t.$$  

We have $g(s) > 0$ and and since $t \mapsto \varphi(\varphi(t,x)) - s$ is bounded, $g(t) < 0$ for $t$ sufficiently small. Hence, by the intermediate value theorem, there is a $t' < s$ such that $g(t') = 0$. Choose $\gamma(n) := \varphi(t',x)$. Then

$$\varphi(\varphi(\gamma(n)), \gamma(n)) = \varphi(s-t', \gamma(n)) = \varphi(s,x) = \gamma(n+1).$$

The following theorem slightly generalizes [14, Theorem 1].

Theorem 3.3. Let $S \subset Y$ be compact. Then the following three conditions are equivalent:

(i) $S$ is an isolated invariant set with respect to $\varphi$.
(ii) For every continuous map $\tau: Y \to \mathbb{R}^+$, $S$ is an isolated invariant set with respect to $\varphi$.  
(iii) There is a number $h > 0$ such that $S$ is an isolated invariant set with respect to $\varphi_h$.

Proof. Assume condition (i) holds and fix a continuous map $\tau: Y \to \mathbb{R}^+$. Choose $N$, an isolating neighborhood for $S$ with respect to $\varphi$. Obviously $\varphi_h(S) \subset S$. Using Lemma 3.2 for every $x \in S$ there is an $x' \in S$ such that $\varphi_h(x') = x$. This means $S \subset \varphi_h(S)$. Hence $S$ is invariant with respect to $\varphi_h$. Let $T := \sup \{ \tau(x) \mid x \in N \} < \infty$. To see that $S$ is an isolated invariant set with respect to $\varphi_h$, consider the map

$$\sigma: N \ni x \mapsto \sup \{ t \in \mathbb{R}_{\geq 0} \mid \varphi([0,t],x) \subset N \} \in [0,\infty].$$

One easily verifies that each $x \in S$ has a compact neighborhood $V_x$ such that $\sigma(V_x) \subset [T,\infty]$. Since $S$ is compact we can choose $M \subset N$, a compact neighborhood of $S$ such that $\sigma(x) \geq T$ for $x \in M$. We will show that $S = \text{Inv}(M, \varphi)$. Obviously we have

$$S = \text{Inv}(N, \varphi) = \text{Inv}(M, \varphi) \subset \text{Inv}(M, \varphi_h).$$

To show the opposite inclusion take $x \in \text{Inv}(M, \varphi_h)$ and let $\gamma: \mathbb{Z} \to Y$ be a solution of $\varphi_h$ through $x$. Let $x_n := \gamma(n)$ and $t_n := \tau(x_n)$. Then

$$x_{n+1} = \varphi_h^{n+1}(x) = \varphi(t_n, \varphi_h^n(x)) = \varphi(t_n, x_n).$$

By definition of $T$, we have $t_n \leq T$. Since $x_n \in M$, we have $\sigma(x_n) \geq T$. It follows that $\varphi([0,t_n],x_n) \subset N$ for all $n$, and consequently $x \in \text{Inv}(N, \varphi) = S$. Thus implication (i) $\implies$ (ii) is proven. Implication (ii) $\implies$ (iii) is obvious because we can always take $\tau$ to be a constant positive function. Implication (iii) $\implies$ (i) is part of Theorem 1 in [14].

The following follows immediately from the implication (iii) $\implies$ (i).

Corollary 3.4 ([13], Lemma 6). Let $N \subset Y$ be an isolating neighborhood for $\varphi_h$. Then

$$\text{Inv}(N, \varphi) = \text{Inv}(N, \varphi_h).$$

Remark 3.5. There is no full analogue of Corollary 3.4 since we cannot replace condition (iii) in the theorem above by the statement
Definition 3.6. \( \text{Let } \phi \text{ be a map. Therefore, we need to consider both situations here.} \)

(\( \text{Def. 3.6} \)) In [6], we propose using algorithms developed for the calculation of Conley indices.  

Definition 3.7. \( \text{Let } P_n \subset N \text{ be compact and } S = \text{Inv}(N, \phi). \text{ A pair } (P_1, P_2) \text{ of compact sets } P_2 \subset P_1 \subset N \text{ is an index pair for } (S, \phi) \text{ in } N \text{ if it has the following properties:} \)

(i) \( \text{If } x \in P_1, t > 0, \phi([0, t], x) \subset N, \text{ then } \phi([0, t], x) \subset P_1; \)
(ii) \( \text{If } x \in P_1, t > 0, \phi(t, x) \notin N, \text{ then } \exists t' \in [0, t] : \phi(t', x) \in P_2, \phi([0, t'], x) \subset N; \)
(iii) \( S \subset \text{int}(P_1 \setminus P_2). \)

For an isolated invariant set \( S \), an index pair exists and the following definition does not depend on a particular choice of the isolating neighborhood \( N \) and the index pair \((P_1, P_2)\).

Definition 3.8. \( \text{Let } (P_1, P_2) \text{ be an index pair for } (S, \phi) \text{ in } N. \text{ The homological Conley index is} \)
\[ \text{CH}_n(S, \phi) := H_n(P_1, P_2). \]

There are several ways of defining the Conley index for maps. Each definition uses a certain kind of index pair and an equivalence relation on a map on this index pair (e.g., [19, 9, 16]). We use the following index pair from [13, Definition 3.1] here and two popular equivalence relations in the subsections 3.2.1 and 3.2.2.

Definition 3.9. \( \text{Let } f : Y \rightarrow Y \text{ be a continuous map. Let } N \subset Y \text{ be compact and } S = \text{Inv}(N, f). \text{ A pair } (P_1, P_2) \text{ of compact sets } P_2 \subset P_1 \subset N \text{ is an index pair for } (S, \phi) \text{ in } N \text{ if} \)

(i) \( f(P_1) \cap N \subset P_1, \)
(ii) \( P_1 \setminus f^{-1}(P_2) \subset P_2, \)
(iii) \( S \subset \text{int}(P_1 \setminus P_2). \)

Let \( H_n \) denote the homology functor with compact supports as defined in [12, Chapter 9]. We use this version of the homology functor because of its good excision properties. However, note that this homology theory coincides with singular homology on compact neighborhood retracts, in particular compact polyhedra and cubical sets.

Let \((P_1, P_2)\) be an index pair for \((S, f)\) in \(N\). Define maps on topological pairs:
\[ f_P : (P_1, P_2) \ni x \mapsto f(x) \in (P_1 \cup (Y \setminus \text{int}N), P_2 \cup (Y \setminus \text{int}N)); \]
\[ i_P : (P_1, P_2) \ni x \mapsto x \in (P_1 \cup (Y \setminus \text{int}N), P_2 \cup (Y \setminus \text{int}N)). \]

The induced homomorphism \( H_n(i_P) \) in relative homology is an isomorphism because of excision (see, [12, Theorem 9.3]). The endomorphism \( I_P = H_n(i_P)^{-1} \circ H_n(f_P) \) on \( H_n(P_1, P_2) \) is called the index map.
In the next two subsections, we use the following notations for an endomorphism $\alpha: V \to V$. Let $\text{dom}(\alpha) = V$ denote the domain of $\alpha$ and $\text{gker}(\alpha)$ the generalized kernel defined by

$$\text{gker}(\alpha) = \{v \in V \mid \exists n \in \mathbb{N}: \alpha^n(v) = 0\}.$$ 

3.2.1. Leray functor. Given an endomorphism $\alpha: V \to V$ of graded modules (e.g. homology groups), the Leray functor $L$ assigns to $\alpha$ an automorphism $L(\alpha)$ as follows. Let

$$\overline{\alpha}: V / \text{gker}(\alpha) \to V / \text{gker}(\alpha), \ [v] \mapsto [\alpha(v)].$$

Then let

$$V' := \cap_{n \in \mathbb{N}} \overline{\alpha}^n(V') \text{ and } L(\alpha): V' \to V', \ [v] \mapsto \overline{\alpha}([v]) = [\alpha(v)].$$

Observe that if $\alpha$ is already an automorphism, then $L(\alpha) = \alpha$. For a detailed definition of this functor, we refer the reader to [16, Section 4]. We define $\alpha \sim \beta$ iff $L(\alpha)$ and $L(\beta)$ are conjugate. Here we let the Conley index be the equivalence class of $L(I_P)$ under this equivalence relation. This definition only depends on $S$ and not on the choice of $P$ or $N$. [16, Theorem 2.6]. For all the index maps $I_P$ considered in this article, the automorphism $L(I_P)$ is conjugate to the identity as is shown later in this section.

The main ingredient for comparing the Conley index of $\phi$ with the Conley index of $\phi_h$ for some $h > 0$ is the existence of a common index pair for both dynamical systems. The following lemma was verified in the proof of [14, Theorem 2].

**Lemma 3.9.** Let $S$ be an isolated invariant set for $\phi_h$ (hence for $\phi$). Then there is a pair $Q' = (Q_1, Q_2)$ of compact sets such that $Q$ is an index pair for $(S, \phi)$ and an index pair for $(S, \phi_h)$.

We also have a slightly more general statement for a time step function $\tau$, but we need an extra assumption about the equality of the invariant sets with respect to $\phi$ and $\phi_\tau$.

**Lemma 3.10.** Let $S$ be an isolated invariant set for $\phi_\tau$ and also for $\phi$. Then there are compact sets $Q_1, Q_2$ and $M$ such that

(i) $Q = (Q_1, Q_2)$ is an index pair for $(S, \phi)$ and for $(S, \phi_\tau)$ in $M$, and

(ii) the index map $I_Q$ is the identity on $H_*(Q_1, Q_2)$.

**Proof.** The proof is similar to the proof of Lemma 3.9 as given in [14, Theorem 2]: The proof of [20, Chapter I, Theorem 5.1] shows the existence of an open neighborhood $V \subset Y$ of $S$ such that $cV$ isolates $S$ and of continuous functions $\kappa, \lambda: V \to [0, \infty)$ such that $S = \kappa^{-1}(0) \cap \lambda^{-1}(0)$, and

(a) If $\kappa(x) > 0, t > 0$ and $\phi(t, x) \in V$, then $\kappa(x) < \kappa(\phi(t, x))$;

(b) if $\lambda(x) > 0, t > 0$ and $\phi(t, x) \in V$, then $\lambda(x) > \lambda(\phi(t, x))$.

For an arbitrary $\zeta > 0$, define subsets of $V$:

$$G(\zeta) := \{x \in V \mid \kappa(x) < \zeta, \lambda(x) < \zeta\},$$

$$H(\zeta) := \{x \in V \mid \kappa(x) \leq \zeta, \lambda(x) \leq \zeta\}.$$ 

Let $T := \max \{ \tau(x) \mid x \in cV \}$. The local compactness of $Y$ can be used to observe that for any open neighborhood $U$ of $S$ there is a $\zeta > 0$ such that $\phi([0, T], H(\zeta)) \subset U$. Applying this observation to $U = \text{int} V$, we conclude the existence of an $\varepsilon > 0$ such that $\phi([0, T], H(\varepsilon)) \subset V$. 
Let $M := H(\varepsilon)$. Applying the observation to $U = G(\varepsilon)$ shows the existence of a $\delta > 0$ such that $\varphi([0, T], H(\delta)) \subset G(\varepsilon)$. Define

$$Q_1 := \{ x \in V | \kappa(x) \leq \varepsilon, \lambda(x) \leq \delta \},$$

$$Q_2 := \{ x \in V | \delta \leq \kappa(x) \leq \varepsilon, \lambda(x) \leq \delta \} \subset Q_1.$$ 

The pair $(Q_1, Q_2)$ is an index pair for $(S, \varphi)$ and for $(S, \varphi_{\tau})$ in $M$. It is straightforward to check the properties in Definitions 3.6 and 3.8, hence (i) is proven. To see (ii), we note that $(Q_1, Q_2) \times [0, 1] \ni (x, s) \mapsto \varphi(s \tau(x), x) \in (Q_1 \cup (Y \setminus \text{int} M), Q_2 \cup (Y \setminus \text{int} M))$

is a homotopy from $i_Q$ to $f_Q$. Hence the index map $I_Q$ is the identity. 

**Theorem 3.11.** Let $S$ be an isolated invariant set for $\varphi$ and $\varphi_{\tau}$. Let $(P_1, P_2)$ be an index pair for $(S, \varphi)$ and $I_P$ its index map. Then there is an isomorphism

$$\text{CH}_*(S, \varphi) \cong \text{dom}L(I_P).$$

**Proof.** Consider a common index pair $(Q_1, Q_2)$ for $(S, \varphi)$ and $(S, \varphi_{\tau})$ in $M$ as in Lemma 3.10.

Then $L(I_P)$ and $L(I_Q)$ are conjugate and therefore

$$\text{dom}L(I_P) \cong \text{dom}L(I_Q) = \text{dom}I_Q = \text{H}_*(Q_1, Q_2).$$

3.2.2. **Shift equivalences.** Another equivalence relation commonly used on the index map $I_P$ is the following one proposed in [19].

**Definition 3.12.** Let $\alpha : V \to V$ and $\beta : W \to W$ be endomorphisms of graded modules. They are **shift equivalent** if there are homomorphisms $r : V \to W$ and $s : W \to V$ such that $\beta \circ r = r \circ \alpha$, $s \circ \beta = \alpha \circ s$, $s \circ r = \alpha^n$ and $r \circ s = \beta^n$ for some $n \in \mathbb{N}$.

**Theorem 3.13.** ([13, Theorem 3.29], [21, Lemma 4.3]) If $P$ and $P'$ are two index pairs around the same isolated invariant set for a discrete dynamical system, then $I_P$ and $I_{P'}$ are shift equivalent.

The following theorem is similar to, but slightly stronger than Theorem 3.11.

**Theorem 3.14.** Let $S$ be an isolated invariant set for $\varphi$ and $\varphi_{\tau}$. Let $(P_1, P_2)$ be an index pair for $(S, \varphi_{\tau})$ and $I_P$ its index map. Then there is an isomorphism

$$\text{CH}_*(S, \varphi) \cong \text{H}_*(P_1, P_2)/\text{gker}(I_P).$$

**Proof.** There is a common index pair $(Q_1, Q_2)$ for $(S, \varphi)$ and $(S, \varphi_{\tau})$ due to Lemma 3.10. Thus, by Theorem 3.13 there are homomorphisms $r, s$ and an $n \in \mathbb{N}$ such that the following diagram commutes.

$$
\begin{array}{ccc}
H_*(P_1, P_2) & \xrightarrow{I_P} & H_*(P_1, P_2) \\
\downarrow r & & \downarrow r \\
H_*(Q_1, Q_2) & \xrightarrow{I_P} & H_*(Q_1, Q_2) \\
\end{array}
$$

The lower right triangle shows that $s$ is injective and $r$ is surjective. This yields

$$\text{CH}_*(S, \varphi) \cong \text{H}_*(Q_1, Q_2) \cong \text{im}(s) = \text{im}(I_P^n) \cong \text{H}_*(P_1, P_2)/\text{ker}(I_P^n).$$

Additionally, $\ker(I_P^n) = \text{gker}(I_P)$ because $I_P^{kn} = I_P^k$ for any $k > 0$. 

□
3.3. Morse decompositions.

Definition 3.15. Let $y \in Y$.

(i) The $\omega$-limit set $\omega(y, \varphi)$ is the set of accumulation points of $\varphi([0, \infty), y)$.

(ii) The $\alpha$-limit set $\alpha(y, \varphi)$ is the set of accumulation points of $\varphi((0, \infty], y)$.

Definition 3.16. Let $f : Y \to Y$ be a discrete dynamical system. Let $y \in Y$ be such that there is a solution $Z \to Y$ through $y$.

(i) The $\omega$-limit set $\omega(y, f)$ is the set of accumulation points of $\{f^k(y) \mid k > 0\}$.

(ii) A point $x \in Y$ is in the $\alpha$-limit set $\alpha(y, f)$ iff there exists a solution $\gamma : \mathbb{Z} \to Y$ with $\gamma(0) = y$ such that $x$ is an accumulation point of $\gamma((0, \infty])$.

Definition 3.17. Given a dynamical system (i.e., a flow $\varphi$ or a map $f$) and an isolating neighborhood $X$ with $S = \text{Inv}(X)$ (denoting $\text{Inv}(X, \varphi)$ or $\text{Inv}(X, f)$, respectively), we call a set of disjoint isolated invariant sets $\{M_p \mid p \in \mathcal{P}\}$ together with an acyclic directed graph $\mathcal{M}G$ with vertices $\mathcal{P}$ a Morse decomposition in $X$ if for every $y \in S$ one of the following holds:

(i) $y \in M_p$ for a certain $p \in \mathcal{P}$; or

(ii) there are $p, q \in \mathcal{P}$ and a path from $p$ to $q$ in $\mathcal{M}G$, such that $\alpha(y) \subset M_p$ and $\omega(y) \subset M_q$.

In order to show that a certain Morse decomposition for $\varphi_t$ is also a Morse decomposition for $\varphi$, we can apply the following two lemmata.

For $A \subset Y$, let

$$\varphi_{[0, \tau]}(A) := \bigcup_{x \in A} \varphi([0, \tau(x)], x) \subset Y.$$ 

Lemma 3.18. Let $\{M_p \mid p \in \mathcal{P}\}$ be a Morse decomposition in $X$ for $\varphi_t$ with isolating neighborhoods $N_p$, i.e., $N_p \cap N_q = \emptyset$ for $p \neq q$ and $M_p = \text{Inv}(N_p, \varphi_t) \subset \text{int}N_p$ for all $p \in \mathcal{P}$. Let $p \in \mathcal{P}$. If

$$(*) \quad \varphi_{[0, \tau]}(N_p) \subset X \text{ and } \varphi_{[0, \tau]}(N_p) \cap N_q = \emptyset \text{ for all } q \in \mathcal{P} \setminus \{p\},$$

then $\text{Inv}(N_p, \varphi) = \text{Inv}(N_p, \varphi_t)$.

Proof. Let $p \in \mathcal{P}$ and $N_p' = \text{cl} \left(X \setminus \bigcup_{q \in \mathcal{P} \setminus \{p\}} N_q\right)$. First we show

$$(5) \quad \text{Inv}(N_p, \varphi) \subset \text{Inv}(N_p', \varphi).$$

Let $\gamma : \mathbb{Z} \to N_p$ be a solution to $\varphi_t$ in $N_p$. Then for each $n \in \mathbb{Z}$: $\varphi([0, \tau(n)], \gamma(n)) \subset N_p'$ by assumption $(*)$ in the lemma. Gluing these pieces together yields a trajectory for $\varphi$ in $N_p'$.

We also show

$$(6) \quad \text{Inv}(N_p', \varphi_t) \subset \text{Inv}(N_p, \varphi_t).$$

Assume there is a point $x \in \text{Inv}(N_p', \varphi_t) \setminus \text{Inv}(N_p, \varphi_t)$. Hence, $x \notin \bigcup_{q \in \mathcal{P}} M_q$. Now either $\alpha(x, \varphi_t)$ or $\omega(x, \varphi_t)$ must lie in some $M_q$ with $q \neq p$ because they cannot lie within the same Morse set. But this means that any solution of $\varphi_t$ through $x$ has to contain points in $\text{int}N_q$, which is disjoint from $N_p'$. We conclude $x \notin \text{Inv}(N_p', \varphi_t)$, a contradiction.

Overall, we get the following inclusions, where the middle one is trivial.

$$\begin{align*}
M_p \subset \text{Inv}(N_p', \varphi) \subset \text{Inv}(N_p', \varphi_t) \subset M_p.
\end{align*}$$

Each set is the same subset of $N_p$. Therefore also $\text{Inv}(N_p', \varphi) = \text{Inv}(N_p, \varphi).$  \qed
Lemma 3.19. Let \( X \subset Y \) be an isolating neighborhood for \( \varphi_t \) (and hence for \( \varphi \)). Let \( \{ M_p \mid p \in P \} \) be a Morse decomposition for \( \varphi_t \) in \( X \) with Morse graph \( MG \) and assume that each \( M_p \) is invariant also for \( \varphi \). Then \( \{ M_p \mid p \in P \} \) is also a Morse decomposition for \( \varphi \) in \( X \) with Morse graph \( MG \).

Proof. We only need to show that the Morse graph \( MG \) is preserved. From here, consider a point \( y \in \text{Inv}(X, \varphi) \subset \text{Inv}(X, \varphi_t) \). There are \( p, q \in P \) and a directed path from \( p \) to \( q \) in \( MG \) such that \( \alpha(y, \varphi_t) \subset M_p \) and \( \omega(y, \varphi_t) \subset M_q \). We show that \( \omega(y, \varphi) \subset M_q \) by contradiction. The analogous statement for the \( \alpha \)-limit is proven similarly.

Let \( N_q \) be an isolating neighborhood for \( \varphi_t \) and therefore for \( \varphi \) around \( M_q \), hence \( \text{Inv}(N_q, \varphi) = \text{Inv}(N_q, \varphi_t) = M_q \subset \text{int}N_q \). We continue analogously to the proof of Theorem 3.3. We define a function

\[
\sigma: N_q \ni x \mapsto \sup \{ t \in \mathbb{R}_{\geq 0} \mid \varphi([0,t],x) \subset N_q \} \in [0,\infty).
\]

Let \( T = \max \tau(N_q) \). Now there is a compact neighborhood \( \tilde{N}_q \) of \( M_q \) such that \( \sigma(x) \geq T \) for all \( x \in \tilde{N}_q \).

Assume that \( \omega(y, \varphi) \) contains a point \( y' \) outside of \( N_q \). We construct a subsequence of \( y \) as follows: Since \( \tilde{N}_q \) is a neighborhood of \( \omega(y, \varphi_t) \) and \( \gamma(N) \cap M_q = \emptyset \), there is an \( \tilde{n} \geq 0 \) such that \( \gamma(\tilde{n}) \in \tilde{N}_q \setminus M_q \). Let \( s > 0 \) be such that \( \gamma(s) \subset \text{int}M_q \). There is an \( s' \geq s \) such that \( \sigma(\varphi(s', y)) = T \), and then \( \varphi([s', s' + T], x) \subset \text{cl}(N_q \setminus \tilde{N}_q) \). Hence, there is an \( n_0 \geq \tilde{n} \) such that \( \gamma(n_0) \in \text{cl}(N_q \setminus \tilde{N}_q) \). Going on like this, we construct a subsequence \( \gamma(n_0), \gamma(n_1), \ldots \) of points in \( \text{cl}(N_q \setminus \tilde{N}_q) \). It has a converging subsequence whose limit is also in \( \omega(y, \varphi_t) \subset M_q \). A contradiction.

Overall, \( \omega(y, \varphi) \subset N_q \) and therefore \( \omega(y, \varphi) \subset \text{Inv}(N_q, \varphi) = M_q \) because \( \omega \)-limits are invariant. \( \square \)

We are ready to show

Theorem 3.20. Let \( X \) be an isolating neighborhood for \( \varphi_t \) for an arbitrary continuous function \( \tau: Y \to \mathbb{R}^+ \) and \( \{ M_p \mid p \in P \} \) a Morse decomposition for \( \varphi_t \) in \( X \) with isolating neighborhoods \( \{ N_p \} \).

Suppose that either

(A) the function \( \tau \) is constant, i.e., \( \tau(x) = h \) for all \( x \in Y \); or

(B) for all \( p \in P \) : \( \varphi_{[0, \tau]}(N_p) \subset X \) and for all \( q \neq p \) holds \( \varphi_{[0, \tau]}(N_p) \cap N_q = \emptyset \).

Then \( \{ M_p \} \) is also a Morse decomposition for \( \varphi \) in \( X \) with the same Morse graph.

Proof. The case \( A \) of a constant time step follows from Corollary 3.4 and Lemma 3.19. The case \( B \) follows from Lemmata 3.18 and 3.19. \( \square \)

Remark 3.21. Criterion \( B \) enables us to verify numerically that the output from the algorithm in Section 2.3 does indeed describe a Morse decomposition: For each \( p \in P \), we construct an enclosure \( \mathcal{Z}(p) \) of all the trajectories \( \varphi([0, \tau(x)], x) \in |N(p)| \), such that \( \mathcal{Z}(p) \subset X \setminus \cup_{q 
eq p} \mathcal{N}(q) \). Our algorithm uses CAPD routines to construct a good enclosure \( \mathcal{Z}(p) \) with

\[
\bigcup_{\xi \in \mathcal{N}(p)} \varphi([0, \max \tau(\xi)], \xi) \subset |\mathcal{Z}(p)|.
\]

In Figure 2(c), each set \( \mathcal{N}(p) \) is shown with a darker collar around it such that together they form \( \mathcal{Z}(p) \). In our example, the algorithm successfully checked \( \mathcal{Z}(p) \subset X \) and \( \mathcal{Z}(p) \cap \mathcal{N}(q) = \emptyset \) for any \( p \neq q \). The additional checks for criterion \( B \) took only about a second.
The following example shows that when [A] is not fulfilled, we need some kind of check that a Morse set $M_p$ for $\varphi_\tau$ is also invariant under $\varphi$.

**Remark 3.22.** Let $S^1 = \{ (\cos \theta, \sin \theta) \mid \theta \in [0, 2\pi) \} \subset \mathbb{R}^2$ be the phase space $X = Y$ and let the flow $\varphi$ be induced by $\dot{\theta} = 1$.

For every point $\theta \in S^1$, its limit sets are $\omega(\theta, \varphi) = S^1$ and $\alpha(\theta, \varphi) = S^1$. The sets $\emptyset$ and $S^1$ are the only subsets invariant for $\varphi$. Define the time step function

$$\tau: S^1 \ni \theta \mapsto \sin \theta + 2\pi \in \mathbb{R}^+,$$

which is well-defined and continuous since it has period $2\pi$. Figure 3 shows a plot of $\tau$ and a trajectory of $\varphi_\tau$.

By definition, $\tau(0) = \tau(\pi) = 2\pi, 2\pi < \tau((0, \pi)) < 3\pi$ and $\pi < \tau((\pi, 2\pi)) < 2\pi$. Additionally $|\tau'(\theta)| < 1$ for $\theta \neq 0, \pi$. These properties suffice to see that for any $0 < \epsilon < \pi$ the subsets $N_1 = [-\epsilon, +\epsilon]$ and $N_2 = [\pi - \epsilon, \pi + \epsilon]$ are isolating neighborhoods for $\varphi_\tau$ with isolated invariant sets $M_1 = \{0\}$ and $M_2 = \{\pi\}$. Whenever $\theta \notin \{0, \pi\}$, then $\alpha(\theta, \varphi_\tau) = M_1$ and $\omega(\theta, \varphi_\tau) = M_2$.

We do therefore get a Morse graph $1 \rightarrow 2$ for $\varphi_\tau$, an attractor-repeller pair. But there is no attractor-repeller pair for $\varphi$ in which both invariant sets are non-empty.\[1\]

The example shows that the assumption that $M_p$ be invariant also for the flow is necessary in Lemma 3.19. Such a function $\tau$ can always be constructed when the flow has a limit cycle $L$. One just has to replace $2\pi$ by the period of the limit cycle and extend $\tau$ from $L$ to the whole phase space $Y$ using Tietze’s extension theorem.

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1R. Vandervorst provided a similar example, independently, to the first author.
4. Final remarks

Using a time step function varying in space gives a lot of flexibility when applying rigorous algorithms for finding Morse decompositions. Sometimes, the step function is the only way to obtain a meaningful Morse decomposition numerically. Additionally, the heuristic presented can be justified more naturally than a specific choice of fixed time step. Up to now, we cannot see from a given equation which strategy is more promising. It would be interesting to find a more systematic approach to find out when to use which strategy.

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DISCRETIZATION STRATEGIES FOR MORSE DECOMPOSITIONS

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