Degeneracy is OK: Logarithmic Regret for Network Revenue Management with Indiscrete Distributions

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We study the classical Network Revenue Management (NRM) problem with accept/reject decisions and $T$ IID arrivals. We consider a distributional form where each arrival must fall under a finite number of possible categories, each with a deterministic resource consumption vector, but a random value distributed continuously over an interval. We develop an online algorithm that achieves $O(\log^2 T)$ regret under this model, with no further assumptions. We develop another online algorithm that achieves an improved $O(\log T)$ regret, with only a second-order growth assumption. To our knowledge, these are the first results achieving logarithmic-level regret in a continuous-distribution NRM model without further “non-degeneracy” assumptions. Our results are achieved via new techniques including: a new method of bounding myopic regret, a “semi-fluid” relaxation of the offline allocation, and an improved bound on the “dual convergence”.

1. Introduction

In the Network Revenue Management (NRM) problem, resources with finite capacities are to be allocated over a finite time horizon of length $T$. During each time step $t = 1, \ldots, T$, a query $t$ arrives, demanding a vector $\tilde{a}_t$ of resources and providing a reward $\tilde{r}_t$. An irrevocable decision must then be made about whether to serve query $t$, in which case $\tilde{a}_t$ would be subtracted from the resources and $\tilde{r}_t$ would be collected. Query $t$ is only feasible to serve if the remaining resources exceed $\tilde{a}_t$ component-wise, and a feasible query $t$ can be judiciously rejected, e.g. if $\tilde{r}_t$ is low relative to $\tilde{a}_t$. The goal is to maximize the total reward collected from serving queries, when the values $(\tilde{r}_t, \tilde{a}_t)$ for each query $t$ are unknown before it arrives but known to be drawn IID across time.

Due to the curse of dimensionality in this problem, a mathematically rich literature has evolved out of developing heuristics and obtaining guarantees on their performance. We consider the stream of literature that analyzes regret, which is the additive loss of an online allocation algorithm compared to an optimal offline allocation that knows all values of $(\tilde{r}_t, \tilde{a}_t)$ beforehand, taken in expectation over the IID query draws (and any further randomness in the algorithm). Generally, regret
is larger for longer time horizons $T$, and this literature is concerned with how the regret grows as a function of $T$ when all other system parameters stay fixed (but the initial resource capacities are also allowed to grow arbitrarily with $T$).

Two types of assumptions are commonly made in the papers that analyze regret in NRM. The first involves having a small number of possible realizations for the values $(\tilde{r}_t, \tilde{a}_t)$, described by a discrete distribution on $N$ points. As $T$ grows, $N$ stays fixed and is treated as a constant in the analysis. However, such an assumption abandons some natural models, e.g. that of $\tilde{r}_t$ being drawn uniformly from $[0,1]$. On the other hand, papers that can capture these continuous distributions require a different set of assumptions, which we will call non-degeneracy. At a high level, these papers assume that the mathematical program being re-solved by the online algorithm over time to make its decisions is always well-behaved, and has a unique optimal solution. However, such assumptions are difficult to intuit or verify, and appear to be motivated primarily by the analysis.

**Our contribution.** We establish logarithmic regret in Network Revenue Management with neither the small-$N$ nor non-degeneracy assumptions. We do make some assumptions on the distribution for $(\tilde{r}_t, \tilde{a}_t)$, but we can show them to be necessary for logarithmic-or-better regret.

To our knowledge, such a result has been missing from the literature, which we now review. Our result builds upon the existing literature that establishes logarithmic-or-better regret in either the small-$N$ or non-degenerate settings, or in the multi-secretary special case.

### 1.1. Regret Results in Network Revenue Management

We begin with Jasin and Kumar (2012), who establish a constant $O(1)$ regret under both the small-$N$ and non-degeneracy assumptions. For NRM literature pre-dating 2012, we refer to the exegesis in Bumpensanti and Wang (2020).

**NRM with small-$N$.** Bumpensanti and Wang (2020) and Vera and Banerjee (2021) were the first to establish $O(1)$ regret for a general NRM problem without any non-degeneracy assumptions. We note that Arlotto and Gurvich (2019) first established $O(1)$ regret in the multi-secretary\(^1\) special case, where all queries demand one unit of a single resource (i.e. $\tilde{a}_t = 1$ w.p. 1). These are all surprising results, in that given a fixed discrete distribution for $(\tilde{r}_t, \tilde{a}_t)$, the regret is upper-bounded by an absolute numerical constant, regardless of how long a time horizon $T$ and how large the resource capacities there are over which regret can be incurred. The intuition is that the further out $t$ is from the end, the harder it is to incur regret, due to the concentration properties of IID draws from a discrete distribution. The total of expected regret over time steps $t = 1, 2, \ldots$ then turns out to be a convergent sum.

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\(^1\)Not to be confused with the use of “secretary” to describe a random arrival order (see e.g. Esfandiari et al. 2017).
**Multi-secretary with general distributions.** A caveat to the aforementioned analysis is that the $O(1)$ constant depends on $N$, and $N$ can be $\infty$, e.g., for continuous distributions. To understand continuous distributions better, Bray (2019) studies the example of $r_i$ being drawn uniformly from $[0,1]$ in the multi-secretary special case. Both Lueker (1998) and Bray (2019) establish an upper bound on regret that grows logarithmically with $T$, and importantly, show this regret rate of $\Theta(\log T)$ to be tight—that is, a constant regret is no longer possible. Recently, Besbes et al. (2022) make further progress on the multi-secretary problem by establishing a notion of complexity for general distributions, which affects regret. A corollary relevant to us is that if the reward distribution is supported on an interval (or many disjoint intervals), over which it has a density that is lower-bounded by a positive constant, then $O(\log^2 T)$ regret can be achieved.

**NRM with general distributions and non-degeneracy.** Although the above papers consider continuous and general reward distributions, it is unclear how they extend beyond the multi-secretary case. Meanwhile, several papers have studied general distributions for NRM under additional assumptions. These assumptions can be classified into the following two categories, which we further discuss in Appendix A:

- **a non-degeneracy condition:** the fluid relaxation of the optimal offline allocation has a unique optimal solution and strict complementary slackness is satisfied.
- **a second-order growth condition:** the Lagrangian dual function of the fluid relaxation is strongly convex.

Though in the literature, these two conditions are always stated together as a whole, we do show that these conditions can indeed be differentiated from each other. In Example 2, we construct an instance where the second-order growth condition is satisfied but the strict complementary slackness of the fluid relaxation is violated.

The non-degeneracy condition is used in the previous literature (Li and Ye 2021, Balseiro et al. 2021, Bray 2022) to guarantee that the binding constraints of the fluid relaxation remain unchanged when the right-hand-side constraint is being replaced by the real-time per period remaining capacities. However, this requires that the real-time per period remaining capacities stay in a neighborhood of the initial per period remaining capacities. In order to guarantee this, the previous literature needs to assume that the initial capacities scale linearly with $T$. This assumption would require: i) the resource with a non-binding constraint in the fluid relaxation to have a buffer capacity that equals $\delta T$ for some fixed constant $\delta > 0$, and; ii) the resource with a binding constraint in the fluid relaxation to have a unique optimal dual variable lower bounded by $\delta$, where the optimal dual variable remains fixed regardless of how large $T$ is since initial capacities scale linearly in $T$ (the regret would scale polynomially in $1/\delta$). However, the square-root law of inventory in practice suggests that the buffer capacity should generally scale as $O(\sqrt{T})$, violating i); and if all resource
constraints are binding, then the optimal dual variables can be non-unique or arbitrarily close to 0. All in all, as highlighted in Bumpensanti and Wang (2020), degeneracy is generally likely to occur in practice, and it remains unknown in the literature whether a logarithmic level regret can be achieved for indiscrete distributions without assuming the non-degeneracy condition and without requiring the initial capacities to scale linearly with $T$.

**Our model: NRM with discrete demands and continuous rewards.** We consider the following structural form of the distribution for $(\tilde{r}_i, \tilde{a}_i)$. First, the demand vector $\tilde{a}_i$ is drawn from a discrete distribution supported on finite possibilities $a_1, \ldots, a_n$. Then, conditional on $\tilde{a}_i = a_j$ for any $j = 1, \ldots, n$, the reward $\tilde{r}_i$ is drawn from a continuous distribution $F_j$ with a lower-bounded density over some interval $[l_j, u_j]$. We develop two algorithms with the following guarantees:

1. $O(\log^2 T)$ regret under our model, with no further assumptions (Section 3);
2. Improved $O(\log T)$ regret, with the second-order growth assumption (Section 4).

In our model, two different indices $j, j'$ can have $a_j = a_{j'}$; moreover, they can have non-overlapping intervals with $l_j < u_j < l_{j'} < u_{j'}$. Thus, with a single resource and $a_j = (1)$ for all $j = 1, \ldots, n$, we can capture the multi-secretary reward distribution with density lower-bounded over disjoint intervals, and our first result achieves a $O(\log^2 T)$ regret matching the corollary from Besbes et al. (2022). We also emphasize that our base model can exhibit all of: degeneracy, lack of strict complementary slackness, and lack of second-order growth. In Section 4, we present a restricted structural form of the query distribution (Assumption 2) which guarantees second-order growth, but still exhibits degeneracy (Example 2). In both cases, our two algorithms are the only known ways to achieve logarithmic or better regret.

Although our model still imposes a structural form on the distribution for $(\tilde{r}_i, \tilde{a}_i)$, it has the following justifications. First, both the aspects of $n$ being finite and the density being lower-bounded are necessary for logarithmic-or-better regret—see Besbes et al. (2022). Second, because we are restricting the distribution instead of the latent mathematical program, our assumption is much easier to intuit—there is a discrete list of $n$ flight itineraries, and a continuous range of customer valuations for each itinerary. Finally, we believe our model is in some sense a natural sandbox for pushing the technical boundaries of the literature, as we now explain in Subsections 1.2 and 1.3.

**Implications for price-based NRM.** It has been known that any regret bound for NRM with general distributions and accept/reject decisions would imply the same regret bound for the price-based NRM problem, where the decision is the price for each query. This correspondence enables our model to have a wider range of application and we elaborate on this correspondence in Appendix B. However, non-degeneracy is still assumed in the previous literature that derives regret bounds for the price-based NRM problem. For example, Jasin (2014) derives a $O(\log T)$ regret bound and Wang and Wang (2022) derives a $O(1)$ revenue gap with respect to the optimal
online policy. However, both Jasin (2014) and Wang and Wang (2022) would require the initial capacities to scale linearly with $T$ and the initial per period capacities to fall into the interior of a set, staying away from the boundary by a fixed positive constant. This assumption is analogous to the non-degeneracy assumption in the accept/reject model.

1.2. $O(\log^2 T)$ Regret with No Further Assumptions

Our $O(\log^2 T)$ result tries to extend the approach of Vera and Banerjee (2021) from discrete distributions, which we now recap. Fix some remaining resource capacities and suppose there are $s$ time steps left. The offline allocation knows for every possible realization $(r, a)$, called a “type”, the remaining number of queries $t$ with $(\tilde{r}_t, \tilde{a}_t) = (r, a)$, denoted by $\tilde{d}_{(r,a)}$. Meanwhile, consider an online algorithm that solves for an optimal fluid packing, which replaces each $\tilde{d}_{(r,a)}$ with its expectation to specify the queries that can be served. The authors compare $\tilde{y}_{(r,a)}$, the number of queries of type $(r, a)$ accepted by the offline, to $\hat{y}_{(r,a)}$, the number of such queries instructed by the fluid packing to accept. Their key argument is that since the number of queries of type $(r, a)$ is growing linearly in $s$ and any two of them are interchangeable, the online algorithm only “makes a mistake” if $\tilde{y}_{(r,a)}$ and $\hat{y}_{(r,a)}$ are distance $\Omega(s)$ apart. This is a highly unlikely event (over the randomness in the offline’s draws of $\tilde{d}_{(r,a)}$) because $\tilde{y}_{(r,a)}$ and $\hat{y}_{(r,a)}$ are generally only $O(\sqrt{s})$ apart, specifically an event with probability $O(e^{-s})$ that when summed over $s$ leads to constant regret.

A new, “semi-fluid” relaxation of offline. Comparing $\tilde{y}_{(r,a)}$ to $\hat{y}_{(r,a)}$ is meaningless under continuous rewards, because there is zero probability of drawing any specific $(r, a)$. To cope, we introduce a new semi-fluid relaxation that amalgamates decisions over queries with the same demand vector. Specifically, we call each $j = 1, \ldots, n$ a type in our model of NRM, and let $\tilde{d}_j$ denote the remaining number of queries with $\tilde{a}_t = a_j$ and $\tilde{r}_t$ drawn from $F_j$. The semi-fluid relaxation knows $\tilde{d}_j$ for all $j$. However, the semi-fluid relaxation differs from the offline allocation in that it collects exactly the “fluid” value

$$\tilde{d}_j \int_{1-\tilde{y}_j/\tilde{d}_j}^1 F_j^{-1}(q)dq$$

when it accepts $\tilde{y}_j$ queries of type $j$. Objective (1) integrates over the $\tilde{y}_j/\tilde{d}_j$ proportion of the $\tilde{d}_j$ type-$j$ queries with the highest rewards, as explained in Section 3. Meanwhile, our algorithm is still based on solving the (fully) fluid packing, whose variables can also be amalgamated into acceptance quantities $\hat{y}_j$ for each type $j$ (and $\tilde{d}_j$ will be replaced $\mathbb{E}[\tilde{d}_j]$, including in the objective (1)). We can then compare the acceptance quantities $\tilde{y}_j$ to $\hat{y}_j$ for the amalgamated types $j$.

Bounding the myopic regret. Unfortunately, because in our model queries of the same type have different rewards, a mistake can be made without requiring $|\tilde{y}_j - \hat{y}_j| = \Omega(s)$. In fact, a mistake only requires drawing a quantile that is above the acceptance proportion $\tilde{y}_j/\tilde{d}_j$ for the semi-fluid
but below $\hat{y}_j / \mathbb{E}[\tilde{d}_j]$ for the fluid (or vice versa), which occurs with probability $|\tilde{y}_j / \tilde{d}_j - \hat{y}_j / \mathbb{E}[\tilde{d}_j]|$. As such, the rough argument that $\hat{y}_j$ and $\tilde{y}_j$ are $O(\sqrt{s})$ apart (and $\mathbb{E}[\tilde{d}_j] = \Omega(s)$) would lead to a mistake probability of $O(1/\sqrt{s})$, and an undesirable overall regret of $O(\sqrt{T})$. Therefore, we instead follow Bray (2019) who argues that for continuous distributions one must quantify the “myopic regret” at each time step (instead of just bounding the probability that it is non-zero). We decompose overall regret in way (see Section 2) such that the regret at a time step can be quantified as

$$V_c^{\text{Semi}}(\tilde{d} + e_j) - \int_{1-\tilde{y}_j/\mathbb{E}[\tilde{d}_j]}^1 F_j^{-1}(q) dq - \hat{y}_j / \mathbb{E}[\tilde{d}_j] \cdot V_{c-a_j}^{\text{Semi}}(\tilde{d}) - \left(1 - \hat{y}_j / \mathbb{E}[\tilde{d}_j]\right) \cdot V_c^{\text{Semi}}(\tilde{d}),$$

(2)

where $c$ denotes the remaining resource capacities, $j$ denotes the type of the current query, and $V_c^{\text{Semi}}(\tilde{d})$ denotes the optimal objective value of the semi-fluid relaxation given remaining resources $c$ and a generic vector $d = (d_1, \ldots, d_n)$ counting the remaining queries of each type.

To upper-bound (2), we take an optimal solution $\hat{y} = (\hat{y}_1, \ldots, \hat{y}_n)$ for $V_c^{\text{Semi}}(\tilde{d} + e_j)$ and modify it into feasible solutions for $V_{c-a_j}^{\text{Semi}}(\tilde{d})$ and $V_c^{\text{Semi}}(\tilde{d})$, which lower-bounds the latter quantities. As long as these feasible solutions can be constructed by modifying only the $j$'th coordinate of $\hat{y}$, we show that (2) is $O((\tilde{y}_j / \tilde{d}_j - \hat{y}_j / \mathbb{E}[\tilde{d}_j])^2)$. Per the earlier discussion, this is roughly $O((1/\sqrt{s})^2) = O(1/s)$, which when summed over $s$ would sufficiently lead to logarithmic regret.

A boundary-attracted algorithm. If the semi-fluid solution $\hat{y}_j$ is close to 0, however, then $\hat{y}_j$ is difficult to modify into a feasible solution for $V_{c-a_j}^{\text{Semi}}(\tilde{d})$—one cannot pack into the reduced capacity $c - a_j$ by only reducing coordinate $\hat{y}_j$. Our strategy is to bypass this boundary situation by tweaking the online algorithm—if there is a risk of this infeasibility, which we show can be identified by the algorithm checking whether $\hat{y}_j = O(\sqrt{s \log s})$, then it always rejects the current type-$j$ query. This effectively sets $\hat{y}_j = 0$ in (2) and avoids having to lower-bound the quantity $V_{c-a_j}^{\text{Semi}}(\tilde{d})$. On the other extreme, our algorithm always accepts if $\hat{y}_j$ is within $O(\sqrt{s \log s})$ of its maximum value $\mathbb{E}[\tilde{d}_j]$. All in all, we use this tweaked version of the fluid re-solving algorithm that is “attracted to boundaries”, which is similar in spirit to the thresholding in Bumpensanti and Wang (2020) and the “conservatism with respect to gaps” in Besbes et al. (2022). We provide a new explanation for it based on our analysis, and it allows us to always upper-bound (2) by $O((\tilde{y}_j / \tilde{d}_j - \hat{y}_j / \mathbb{E}[\tilde{d}_j])^2) = O(1/s)$ while sacrificing only a log-factor—ultimately achieving $O(\log^2 T)$ regret.

Lipschitz property for semi-fluid convex program. Finally, it should not be taken for granted that $\hat{y}_j$ and $\tilde{y}_j$ are nearby, an intuition we have been frequently using. Indeed, they correspond to optimal solutions of mathematical programs with different objective functions—the fluid problem replaces $\tilde{d}_j$ with $\mathbb{E}[\tilde{d}_j]$ in (1)—and a simple example (Mangasarian and Shiau 1987, Remark 2.7) reveals that optimal solution sets are highly sensitive to small perturbations in the objective. Nonetheless, we extend (Lemma 3) the Lipschitz analysis of Mangasarian and Shiau
(1987) to show that the specific objective function (1) is well-behaved. We also note that due to degeneracy, it is necessary for the Lipschitz property to be of the form “given any optimal solution $\hat{y} = (\hat{y}_1, \ldots, \hat{y}_n)$ to the fluid, there exists a nearby optimal solution $\tilde{y}$ to the semi-fluid”. General perturbation analysis results for convex programs (Bonnans and Shapiro 2013), which try to argue that all optimal solutions $\tilde{y}$ are nearby, do not apply in our setting with degeneracy.

1.3. $O(\log T)$ Regret with Second-order Growth Assumption

Based on the structural form introduced previously, we further assume that the density is also upper bounded, conditional on $\tilde{a}_t = a_j$, and it holds that $l_j = 0$, for any $j = 1, \ldots, n$. Note that the further assumptions we make are to guarantee a second-order growth condition over the Lagrangian dual function of the ex-ante relaxation, which requires the dual function to be strongly convex. As discussed earlier, the second-order growth condition, together with the non-degeneracy assumption, lead to the $O(\log T \log \log T)$ regret bound in Li and Ye (2021) and the $O(\log T)$ regret bound in Balseiro et al. (2021) and Bray (2022). In comparison to the aforementioned papers, our contribution here is to derive a $O(\log T)$ regret bound, without the non-degeneracy assumption, though we still need the second-order growth condition. In what follows, we first illustrate how the two conditions are utilized in the previous literature to derive the logarithmic regret bound, and we then explain how our analysis improves upon them to relax the non-degeneracy assumption.

Analysis in the previous literature. Note that after assuming the strict complementary slackness condition is satisfied by the ex-ante relaxation, there exists a neighborhood of the initial average capacities per period such that the optimal basis (binding constraints) for the ex-ante relaxation remains unchanged even if the average capacities are perturbed, as long as the perturbed average capacities still stay in this neighborhood. Then, one can establish a connection between the average remaining capacities and the optimal dual variables of the ex-ante relaxation, for each period, as long as the average remaining capacities stay in this neighborhood. Such a connection reveals that the dual variables corresponding to the binding constraints will behave as a martingale, while the concentration can be used to bound the regret with the second-order growth condition of the dual function. Such an idea is first developed in Jasin and Kumar (2012) with a small-$N$ assumption, and then extended in Li and Ye (2021) and Balseiro et al. (2021) for general distributions. However, we note that all three aforementioned papers use the ex-ante relaxation as the upper bound of the optimal offline allocation to do the regret analysis, which is the key reason why the non-degeneracy condition is inevitable to carry out their analysis.

A tighter relaxation of the offline allocation. Our improvement comes from using a tighter relaxation than the ex-ante relaxation as an upper bound of the offline allocation. Our relaxation is that for each sample path, we relax the integral decision of the offline allocation to be fractional,
and we take an expectation over the sample path. Such an LP relaxation of the offline allocation has been derived in Bumpensanti and Wang (2020) and Vera and Banerjee (2021) in the discrete setting and we derive it here for general distributions. Then, by following a myopic regret approach described in Section 2.1, we are able to bound the regret incurred at each period by the variance of the dual variable of the LP relaxation of the offline allocation, no matter what the remaining capacities are. Note that we do not need to consider the optimal basis of the ex-ante relaxation to bound the myopic regret. This is the key distinction between our approach and the martingale-based approach in Li and Ye (2021), Balseiro et al. (2021), Bray (2022), which would require a non-degeneracy assumption to guarantee the optimal basis remains fixed in their analysis.

**An improvement on the dual convergence bound.** Another improvement we make is regarding the variance of the dual variable of the LP relaxation of the offline allocation, which is referred to as “dual convergence” in Li and Ye (2021) and requires the second-order growth condition of the Lagrangian dual function of the ex-ante relaxation. Note that if one assumes the distribution to be known, then the martingale-based approach in Li and Ye (2021) would not require a bound on the “dual convergence”. Instead, as shown in Theorem 1 of Balseiro et al. (2021), the martingale property itself is enough to prove the regret bound. The “dual convergence” is needed in Li and Ye (2021) because they allow the distribution to be unknown and they use the sample average problem to estimate the dual variable, where the data comes from the past periods, and the “dual convergence” is needed to bound the estimation error. To be specific, when there are \( s \) data points, Li and Ye (2021) prove the “dual convergence” bound to be at the order of \( O\left(\frac{\log \log s}{s}\right) \), which we improve to \( O\left(\frac{1}{s}\right) \). To obtain this improvement, we utilize both ways of splitting the whole space into a set of small cubes with exponentially increasing edge lengths in Huber (1967) and Li and Ye (2021) to give a probability bound on difference scenarios. To illustrate the main idea, we denote by \( \hat{\mu} \) the dual variable of the sample average problem and \( \tilde{\mu} \) the dual variable of the ex-ante relaxation. Then, we apply the approach in Huber (1967) to obtain a bound on \( P(\|\hat{\mu} - \tilde{\mu}\| \geq \epsilon) \), which is derived from Chebyshev’s inequality. However, we note that when \( \epsilon > \sqrt{\frac{1}{s}} \), the approach in Li and Ye (2021), which is based on the Hoeffding’s inequality, would give us a tighter probability bound. Therefore, by applying different ways to bound \( P(\|\hat{\mu} - \tilde{\mu}\| \geq \epsilon) \) for different ranges of \( \epsilon \), we get a tighter bound on the “dual convergence” \( \mathbb{E}[\|\hat{\mu} - \tilde{\mu}\|^2] \).

We do note that the “dual convergence” bound is used in different ways between our analysis vs. Li and Ye (2021). The data points for “dual convergence” in Li and Ye (2021) comes from past periods. By contrast, the “data points” for “dual convergence” in our analysis comes from the future periods. This is because our online decision is made based on the ex-ante relaxation, while our benchmark, the LP relaxation of the offline allocation, makes the decision based on each sample path of future periods. The myopic regret caused by this distinction is shown to be bounded by the
variance of the dual variable, where the randomness comes from the sample path of future periods. Though the future sample path is convoluted, we only use its distribution to give a bound and our algorithm does not require any knowledge of the realization.

**Discussion on the second-order growth condition.** We do acknowledge that the second-order growth condition that is assumed in Section 4 is somewhat stronger than the one in existing literature (Li and Ye 2021, Balseiro et al. 2021, Bray 2022). To be specific, the second-order growth condition assumed in Section 4 holds for the Lagrangian dual function given any remaining average capacities per period, while the condition in the existing literature holds given remaining average capacities per period belonging to a neighborhood of the initial average capacities per period. However, the second-order growth condition in Section 4 is a consequence of our problem formulation instead of a primitive assumption. Moreover, since it is assumed that the ex-ante relaxation admits a unique optimal dual variable in Li and Ye (2021), Balseiro et al. (2021), Bray (2022), their second-order growth conditions are stated as the strong-convexity of the Lagrangian dual function. In contrast, the second-order growth condition in our setting is stated as the strong-convexity after projecting every variable into the subspace that is spanned by the set of possible query sizes and we do not require the uniqueness of the optimal dual variable.

1.4. Further Related Work

The network revenue management (NRM) problem has been extensively studied in the literature and one main topic is to develop near-optimal policies with strong theoretical guarantees. One common way is to derive the policy from the optimal solution of the ex-ante relaxation. To be specific, Talluri and Van Ryzin (1998) propose a static bid-price policy based on the optimal dual variable of the ex-ante relaxation and proves that the regret bound is $O(\sqrt{T})$. Then, a dynamic update of the bid-price is considered in the literature. Subsequently, Reiman and Wang (2008) shows that by re-solving the ex-ante relaxation once to update the bid-price, one can obtain an improved regret bound $o(\sqrt{T})$. Then, Jasin and Kumar (2012) shows that under a non-degeneracy condition for the ex-ante relaxation, a policy which re-solves the ex-ante relaxation at each time period will lead to an $O(1)$ regret. The relationship between the performances of the control policies and the number of times of re-solving the ex-ante relaxation is further discussed in their later paper (Jasin and Kumar 2013). More recently, Bumpensanti and Wang (2020) proposes an infrequent re-solving policy and shows a regret bound of $O(1)$ without the “non-degeneracy” assumption. This has been extended by Balseiro and Xia (2022) to fair allocation problems. With a different approach, Vera and Banerjee (2021) proves the same $O(1)$ upper bound for the NRM problem and their approach is further generalized in series of papers (e.g. Freund and Banerjee (2019), Vera et al. (2021), Freund and Zhao (2022)). Recent studies on the NRM problem includes variants
such as the reusable resource setting (Baek and Ma 2022), unknown distribution setting (Li et al. 2020, Balseiro et al. 2022) and imperfect distribution knowledge setting under a non-stationary environment (Jiang et al. 2020).

Another problem that is closely related to the NRM problem is called the online packing problem, where a more general formulation is studied and less distribution knowledge is assumed. The packing problem covers a wide range of applications, including secretary problem (Ferguson 1989, Arlotto and Gurvich 2019), online knapsack problem (Arlotto and Xie 2020, Jiang and Zhang 2020), resource allocation problem (Asadpour et al. 2020), network routing problem (Buchbinder and Naor 2009), matching problem (Mehta et al. 2007), etc. The problem is usually studied under either a stochastic model where the reward and size of each query is drawn independently from an unknown distribution $P$, or a more general the random permutation model where the queries arrive in a random order (Molinaro and Ravi 2014, Agrawal et al. 2014, Kesselheim et al. 2014, Gupta and Molinaro 2014).

2. Problem Formulation and Our Approach

We consider an online resource allocation problem, where there are $m$ resources and each resource $i \in [m]$ has an initial fractional capacity $C_i \in \mathbb{R}_{\geq 0}$. There are $T$ discrete time periods and at each period $t \in [T]$, one query arrives, denoted by query $t$. Each query $t$ has a random size $\tilde{a}_t = (\tilde{a}_{t,1}, \ldots, \tilde{a}_{t,m}) \in \mathbb{R}^m_{\geq 0}$, where $\tilde{a}_{t,i}$ denotes how much resource $i$ will be consumed if query $t$ is served, for all $i \in [m]$, and a random reward $\tilde{r}_t \in \mathbb{R}_{\geq 0}$ that denotes how much reward can be collected by serving query $t$. We assume that the value of $(\tilde{r}_t, \tilde{a}_t)$ for each $t \in [T]$ is drawn independently from an identical distribution denoted by $F(\cdot)$. We suppose that the queries are of finite types, i.e., for each $t \in [T]$, $\tilde{a}_t$ is supported on a finite set $A = \{a_1, \ldots, a_n\}$. We call the situation where $\tilde{a}_t$ is realized as $a_j$ as query $t$ being of type $j$ and we denote by $p_j = P(\tilde{a} = a_j)$, for each $j \in [n]$.

After query $t$ arrives and the value of $(\tilde{r}_t, \tilde{a}_t)$ is revealed, the decision maker has to decide immediately and irrevocably whether or not to serve query $t$. Note that query $t$ can only be served if for every resource $i$ its remaining capacity $C_i$ at least $\tilde{a}_{t,i}$. The goal of the decision maker is to maximize the total collected reward subject to the resource capacity constraint.

Any online policy $\pi$ for the decision maker is specified by a set of decision variables $\{\tilde{x}_t^\pi\}_{t \in [T]}$, where $\tilde{x}_t^\pi$ is a binary variable and denotes whether query $t$ is served, for all $t \in [T]$. Note that $\tilde{x}_t^\pi$ can be stochastic if $\pi$ is a randomized policy. Any policy $\pi$ is feasible if for all $t \in [T]$, $\tilde{x}_t^\pi$ depends only on $F(\cdot)$ and $\{(\tilde{r}_1, \tilde{a}_1), \ldots, (\tilde{r}_t, \tilde{a}_t)\}$, and the following capacity constraint is satisfied:

$$\sum_{t=1}^T \tilde{a}_{t,i} \cdot \tilde{x}_t^\pi \leq C_i, \forall i \in [m]. \quad (3)$$
The total collected value of policy $\pi$ is given by $V^\pi(I) = \sum_{t=1}^T \tilde{r}_t \cdot \tilde{x}_t^\pi$, where $I = \{(\tilde{r}_t, \tilde{a}_t)\}_{t=1}^T$ denotes the problem instance.

The benchmark is the prophet, which is an offline decision maker that is aware of the value of $(\tilde{r}_t, \tilde{a}_t)$ for all $t \in [T]$ and always makes the optimal decision in hindsight. We denote by $\{\tilde{x}_{t}^{\text{off}}\}_{t=1}^T$ the offline decision of the prophet, which is an optimal solution to the following offline problem:

$$V^{\text{off}}(I) = \max_{\tilde{a}} \sum_{t=1}^T \tilde{r}_t \cdot \tilde{x}_{t}^{\pi}$$

subject to

$$\sum_{t=1}^T \tilde{a}_{t,i} \cdot x_t \leq C_i, \quad \forall i \in [m]$$

$$x_t \in \{0, 1\}, \quad \forall t \in [T].$$

For any feasible online policy $\pi$, we use regret to measure its performance, which is defined as follows:

$$\text{Regret}(\pi) := \mathbb{E}_{I \sim F}[V^{\text{off}}(I)] - \mathbb{E}_{I \sim F}[V^\pi(I)]$$

where $I = \{(\tilde{r}_t, \tilde{a}_t)\}_{t=1}^T \sim F$ denotes that $(\tilde{r}_t, \tilde{a}_t)$ follows distribution $F(\cdot)$ independently for each $t \in [T]$. In what follows, we describe our general approach to upper bound the regret defined in (5), and discuss how our approach implies online policies under various settings.

### 2.1. General Description of Our Approach

We now give a general description of our approach. We denote by $c = (c_1, \ldots, c_m) \in \mathbb{R}^m$ any vector of remaining capacities of the resources at the beginning of a period $t$. Then, on problem instance $I_t = \{(\tilde{r}_t, \tilde{a}_t)\}$, we denote by $\check{V}_c(I_t)$ a relaxation of the total reward collected by the prophet from period $t$ up to period $T$, given the remaining capacity $c$, where the decision variable $x_t \in \{0, 1\}$ is relaxed into $x_\tau \in [0, 1]$ for $\tau = t, \ldots, T$. We specify various formulations of the relaxation $\check{V}_c(I_t)$ to deal with various settings in the following sections. Then, the regret of any online policy $\pi$ can be upper bounded by the gap between $\mathbb{E}_{I_t \sim F}[\check{V}_C(I_t)]$, where $C = (C_1, \ldots, C_m)$ is a vector of initial capacity for all resources, and $\mathbb{E}_{\pi,I_t \sim F}[V^\pi(I_t)]$, i.e.,

$$\text{Regret}(\pi) \leq \mathbb{E}_{I_t \sim F}[\check{V}_C(I_t)] - \mathbb{E}_{\pi,I_t \sim F}[V^\pi(I_t)].$$

Our approach relies on the following decomposition of the upper bound in (6). For each $t \in [T]$, we denote by $\tilde{c}_t^\pi = (\tilde{c}_{t,1}^\pi, \ldots, \tilde{c}_{t,m}^\pi) \in \mathbb{R}^m$ the remaining capacities at the beginning of period $t$ during the execution of the policy $\pi$. Note that $\tilde{c}_t^\pi$ is random for each $t \in [T]$, where the randomness comes from the randomness in the problem instance $I$ and any randomness in the policy $\pi$. Then, the term $\check{V}_C(I_t)$ can be telescoped as follows by noting that $\tilde{c}_1^\pi = C$ and $\check{V}_c(I_{T+1}) = 0$ for every $c$:

$$\check{V}_C(I_t) = \check{V}_{\tilde{c}_t^\pi}(I_t) - \check{V}_{\tilde{c}_{t+1}^\pi}(I_{t+1}) = \sum_{t=1}^T \left(\check{V}_{\tilde{c}_t^\pi}(I_t) - \check{V}_{\tilde{c}_{t+1}^\pi}(I_{t+1})\right).$$
Thus, the regret upper bound (6) can be decomposed as:

\[
\mathbb{E}_{I_1 \sim F}[\tilde{V}_C(I_1)] - \mathbb{E}_{\pi, I_1 \sim F}[V^\pi(I_1)] = \mathbb{E}_{\pi, I_1 \sim F} \left[ \sum_{t=1}^{T} \left( \tilde{V}_{c_t}^\pi(I_t) - \tilde{V}_{c_{t+1}}^\pi(I_{t+1}) - \tilde{r}_t \cdot \tilde{x}_t^\pi \right) \right]
\]

\[
= \sum_{t=1}^{T} \mathbb{E}_{\pi, I_t \sim F} \left[ \tilde{V}_{c_t}^\pi(I_t) - \tilde{V}_{c_{t+1}}^\pi(I_{t+1}) - \tilde{r}_t \cdot \tilde{x}_t^\pi \right]
\]

\[
= \sum_{t=1}^{T} \mathbb{E}_{\pi, I_t \sim F} \left[ \tilde{V}_{c_t}^\pi(I_t) - \tilde{V}_{c_{t+1}-\tilde{a}_t}^\pi(I_{t+1}) - \tilde{r}_t \cdot \tilde{x}_t^\pi \right]
\]

where the third equality follows from the identity that \( \tilde{c}_{t+1}^\pi = \tilde{c}_t^\pi - \tilde{a}_t \cdot \tilde{x}_t^\pi \). We proceed to analyze the term for each \( t \in [T] \) in the above summation. For each \( c \geq 0 \), we now denote by

\[
\text{Myopic}_t(\pi, c) = \mathbb{E}_{\pi, I_t} \left[ \tilde{V}_c(I_t) - \tilde{V}_{c-\tilde{a}_t}z_t^\pi(I_{t+1}) - \tilde{r}_t \cdot \tilde{x}_t^\pi \right].
\]  

(8)

It is clear that in order to upper bound Regret(\( \pi \)), it is sufficient to upper bound \( \text{Myopic}_t(\pi, c) \) for each \( t \in [T] \) and each \( c \geq 0 \). We summarize the above arguments in the following lemma.

**Lemma 1.** For any feasible online policy \( \pi \), the regret is upper bounded by

\[
\text{Regret}(\pi) \leq \sum_{t=1}^{T} \mathbb{E}_{c_t} [\text{Myopic}_t(\pi, \tilde{c}_t^\pi)]
\]

where the myopic term \( \text{Myopic}_t(\pi, \tilde{c}_t^\pi) \) is defined in (8).

We now motivate our policy \( \pi \) such that the myopic term \( \text{Myopic}_t(\pi, c) \) can be minimized for each \( c \). Now suppose that the online decision maker is allowed to “foresee” the sample path \( I_t \) and we denote by

\[
M_{c, \tilde{a}_t}(I_{t+1}) = \tilde{V}_c(I_{t+1}) - \tilde{V}_{c-\tilde{a}_t}(I_{t+1})
\]

the marginal increase for the relaxation \( \tilde{V} \) to have an extra \( \tilde{a}_t \) resources from period \( t+1 \) to \( T \). Clearly, in order to minimize \( \text{Myopic}_t(\pi, c) \) in (8), we set \( \tilde{x}_t^\pi = 1 \) if and only if

\[
\tilde{V}_c(I_t) - \tilde{V}_{c-\tilde{a}_t}(I_{t+1}) - \tilde{r}_t \leq \tilde{V}_c(I_t) - \tilde{V}_c(I_{t+1})
\]

which implies that

\[
\tilde{x}_t^\pi = \begin{cases} 
1, & \text{if } \tilde{r}_t \geq M_{c, \tilde{a}_t}(I_{t+1}) \\
0, & \text{if } \tilde{r}_t < M_{c, \tilde{a}_t}(I_{t+1}). 
\end{cases}
\]

However, note that in order for \( \pi \) to be feasible, \( \tilde{x}_t^\pi \) must be independent of \( I_{t+1} \). Therefore, instead of comparing \( \tilde{r}_t \) to the marginal increase \( M_{c, \tilde{a}_t}(I_{t+1}) \), we compare \( \tilde{r}_t \) to an estimator \( \tilde{M}_{c, \tilde{a}_t} \) that is independent of \( I_{t+1} \). Our policy is formalized in Algorithm 1, which takes as an input an exogenous estimator \( \tilde{M} \) that we further specify in the following sections on different settings.
Algorithm 1 $\hat{M}$-estimator policy ($\pi_{\hat{M}}$)

1: Input: an estimator $\hat{M}$.
2: Initialize the initial capacities $c_1 = C$.
3: for $t = 1, \ldots, T$ do
4: Observe the value of $(\hat{r}_t, \hat{a}_t)$ and obtain the value of $\hat{M}_{\hat{c}_t, \hat{a}_t}$.
5: if $\hat{r}_t \geq \hat{M}_{\hat{c}_t, \hat{a}_t}$ and $c_t \geq \hat{a}_t$, then we set $x_t = 1$ and update $c_{t+1} = c_t - \hat{a}_t$;
6: otherwise, set $x_t = 0$ and update $c_{t+1} = c_t$.
7: end for
8: Output: online decisions $x = (x_1, \ldots, x_T)$.

3. Policy with Log-squared Regret

In this section, we derive a log-squared bound for (6) under the following assumption over the distribution $F(\cdot)$.

Assumption 1. We assume that for each $j \in [n]$, conditional on $\hat{a}_t$ being realized as any $a_j \in A$, the reward distribution of $\hat{r}$ is supported on the interval $[l_j, u_j]$ with a density function $f(\cdot|a_j)$, where $u_j \geq l_j \geq 0$, and it satisfies $f(r|a_j) \geq \alpha$, for a constant $\alpha > 0$, for any $r \in [l_j, u_j]$.

For notation simplicity, we denote by $F_j(\cdot) = F(\cdot|a_j)$ and $f_j(\cdot) = f(\cdot|a_j)$. Note that in the above Assumption 1, we allow $l_j = u_j$ for a type $j$, i.e., the reward distribution for type $j$ query is a point mass. In this case, we let the density be $f_j(r) = \infty$ for $r = l_j = u_j$ and any constant $\alpha$ would satisfy $f_j(r) \geq \alpha$ for $r \in [l_j, u_j]$.

In what follows, we first specify the relaxation $\tilde{V}$ that will be used in this section, and then we specify the $\hat{M}$-estimator for our algorithm and derive the corresponding regret bound.

3.1. Semi-fluid Relaxation

We now specify the semi-fluid relaxation $\tilde{V}$ that will be used in this section. For each $j \in [n]$, we let $d_j$ denote a generic non-negative integer that should be interpreted as the number of type $j$ query arrivals remaining, irregardless of the current time period. Meanwhile, we let $\tilde{d}_{j,t}$ be the random variable for the number of type $j$ query arrivals from period $t$ to period $T$, i.e., the number of times that $\tilde{a}_\tau = a_j$ for $\tau = t, \ldots, T$. We introduce $\tilde{V}$ is as follows, for a fixed $d = (d_1, \ldots, d_n) \in \mathbb{R}^n$:

$$\tilde{V}^{\text{semi}}(d) = \max_{x} \sum_{j=1}^{n} d_j \cdot \mathbb{E}_{r \sim F_j} [r \cdot x_j(r)]$$

$$\text{s.t. } \sum_{j=1}^{n} d_j \cdot a_{j,i} \cdot \mathbb{E}_{r \sim F_j} [x_j(r)] \leq c_i, \forall i \in [m]$$

$$x_j(r) \in [0, 1], \forall j \in [n], \forall r \in [l_j, u_j]$$
In the following lemma, we show that the formulation of $\tilde{V}$ introduced in (9) implies an upper bound of the offline optimum $V^{\text{off}}$ in (4).

**Lemma 2.** It holds that $E_t[\tilde{V}^\text{Semi}_C(\tilde{d}_1)] \geq E_t[V^{\text{off}}(I)]$, where $\tilde{d}_1 = (\tilde{d}_{1,1}, \ldots, \tilde{d}_{n,1})$ depends on the sample path $I$.

In order to see that $\tilde{V}^\text{Semi}_C(\tilde{d}_1)$ is an upper bound, we first fix the type arrivals of the queries that is implied by the sample path $I$. We then take an ex-ante relaxation over the reward distribution for each type.

**Comparison with other relaxations.** There are also other relaxations of the prophet (4) existing in the literature and we now compare $\tilde{V}^{\text{Semi}}_C(d)$ (9) with them.

The most natural relaxation of the prophet (4) is an LP relaxation, which is defined for each $t \in [T]$, any $c \geq 0$, and any sample path $I$.

$$\tilde{V}^\text{Off}_c(I_t) = \max_x \sum_{\tau=t}^T \tilde{r}_\tau \cdot x_\tau$$

s.t. $\sum_{\tau=t}^T \tilde{a}_{\tau,i} \cdot x_\tau \leq c_i, \forall i \in [m]$

$x_\tau \in [0,1], \forall \tau = t, \ldots, T$

The only difference between the formulation of $V^{\text{off}}(I)$ and $\tilde{V}^\text{Off}_C(I_1)$ is that the integral decision variables of $V^{\text{off}}(I)$ are relaxed to be fractional in $\tilde{V}^\text{Off}_C(I_1)$.

Another common relaxation in the literature is the so-called ex-ante relaxation, which can be obtained by taking expectation over $\tilde{d}_t$ in the formulation of $\tilde{V}^\text{Off}_C(\tilde{d}_t)$. For any $t \in [T]$ and any $c$, we denote by $\tilde{V}^\text{Fld}_{t,c}$ the ex-ante relaxation with a formulation given as follows:

$$\tilde{V}^\text{Fld}_{t,c} = \max_x \sum_{j=1}^n p_j \cdot s \cdot E_{r \sim F_j} [r \cdot x_j(r)]$$

s.t. $\sum_{j=1}^n p_j \cdot a_{j,i} \cdot E_{r \sim F_j} [x_j(r)] \leq c_i, \forall i \in [m]$

$x_j(r) \in [0,1], \forall j \in [n], \forall r \in [l_j, u_j]$

where we denote by $s = T - t + 1$ for notation brevity.

**3.2. Policy and Regret Analysis**

We now develop the estimator that will be used in Algorithm 1 and analyze the regret bound. It is easy to see that the optimal solution of (9) preserves a “threshold” property, i.e. there exists a set of thresholds $\{\kappa_j\}_{j=1}^n$ such that it is optimal to set $x^*_j(r) = 1$ if and only if $r \geq \kappa_j$ and $x^*_j(r) = 0$
if and only if \( r < \kappa_j \), for any \( j \in [n] \). Therefore, \( \bar{V}^{\text{Semi}}_c(d) \) in (9) can be re-written into the following formulation:

\[
\bar{V}^{\text{Semi}}_c(d) = \max_{q} \sum_{j=1}^{n} d_j \cdot \int_{q=1-q_j}^{1} F^{-1}_j(q) dq \\
\text{s.t. } \sum_{j=1}^{n} d_j \cdot a_{j,i} \cdot q_j \leq c_i, \quad \forall i \in [m] \\
q_j \in [0,1], \quad \forall j \in [n].
\] (12)

Here, the decision variable \( q_j \) can be interpreted as the probability of serving type \( j \) query, for each \( j \in [n] \). Denote by \( \{\tilde{q}_j^*\} \) one optimal solution to \( \bar{V}^{\text{Semi}}_c(d) \). Since \( \{\tilde{q}_j^*\} \) depends on the sample path \( I_t \), clearly, one cannot directly use \( \{\tilde{q}_j^*\} \) to derive a feasible online policy that is “agnostic” about \( I_{t+1} \). Therefore, we will consider using the optimal solution of the ex-ante problem (11) to “approximate” \( \{\tilde{q}_j^*\} \). It is clear to see that the optimal solution of the ex-ante problem (11) also preserves a threshold property and \( \bar{V}^{\text{Fld}}_{t,c} \) (11) can be re-written into the following formulation:

\[
\bar{V}^{\text{Fld}}_{t,c} = \max_{q} \sum_{j=1}^{n} p_j \cdot s \cdot \int_{q=1-q_j}^{1} F^{-1}_j(q) dq \\
\text{s.t. } \sum_{j=1}^{n} p_j \cdot s \cdot a_{j,i} \cdot q_j \leq c_i, \quad \forall i \in [m] \\
q_j \in [0,1], \quad \forall j \in [n]
\] (13)

Note that the formulation of \( \bar{V}^{\text{Fld}}_{t,c} \) deviates from the formulation of \( \bar{V}^{\text{Semi}}_c(d) \) only in that the random variable \( \tilde{d}_{j,t} \) is changed into its expectation \( p_j \cdot s \). We can bound how this change of parameter would result in a change of the expected reward that we can gain from each type of query in the optimization problem (12) and (13). Our analysis generalizes the Lipschitz analysis in Mangasarian and Shiau (1987) from linear programming to a general convex optimization problem. Our argument is formulated in the following lemma, where the proof is relegated to Appendix D.

**Lemma 3.** There exists a constant \( \mu \) such that for any \( c \geq 0 \) and any optimal solution \( \hat{q}^* = (\hat{q}_j^*)_{j=1}^{n} \) to (13), it holds that

\[
\|\hat{q}^* - \tilde{q}^*\|_{\infty} \leq \max_{j \in [n]} \left\{ \frac{\mu}{p_j} \right\} \cdot \max_{j \in [n]} \left\{ |p_j - d_j/s| \right\} + \max_{j \in [n]} \left\{ \left| \frac{d_j}{p_j \cdot s} - 1 \right| \right\}, \quad \forall j \in [n]
\]

for any sample path \( I_t \), where \( s = T - t + 1 \), \( \tilde{q}^* = (\tilde{q}_j^*)_{j=1}^{n} \) denotes one optimal solution to (12).

Note that the constant \( \mu \) in Lemma 3 depends only on the matrix \( A = (a_{i,j})_{i \in [m], j \in [n]} \in \mathbb{R}^{m \times n} \), and is thus common for any \( s, c \) and any parameters \( \{l_j, u_j, \alpha\}_{j=1}^{n} \) given in Assumption 1. The exact formulation for the constant \( \mu \) is given in (75) in Appendix D. Then, an upper bound for \( \mu \) can be obtained as in the following lemma, where the proof is relegated to Appendix D.
Lemma 4. Denote by \( \bar{a} = \max_{j \in [n]} \left\{ \sqrt{2 + \sum_{i \in [n]} a_{i,j}^2} \right\} \). Then, for the constant \( \mu \) specified in Lemma 3, it holds that
\[
\mu \leq (m + 2n) \cdot \bar{a}^n. \tag{14}
\]
Lemma 3 implies the existence of a constant \( \kappa_1 \) such that
\[
\max_{j \in [n]} |\hat{q}_j^* - \tilde{q}_j^*| \leq \kappa_1 \cdot \max_{j \in [n]} \{|d_j/s - p_j|\}
\]
which will drive our regret analysis. The formal policy is given in Algorithm 2, which also depends on the constant \( \kappa_1 \) (an upper bound on \( \kappa_1 \) suffices). From (14), we can set
\[
\kappa_1 = \max_{j \in [n]} \left\{ \frac{1}{p_j} \right\} \cdot (m + 2n) \cdot \bar{a}^n. \tag{15}
\]
We now provide the regret analysis.

Algorithm 2 Algorithm achieving \( O(\log^2 T) \) Regret

1: Input: the remaining inventory \( c \), a constant \( \kappa_1 \) given in (15), and the type of query \( t \), denoted by \( j_t \).
2: Obtain \( \{\hat{q}_{j_t,t}^*\} \) by solving the optimization problem (13).
3: if \( \hat{q}_{j_t,t}^* \geq 1 - 3\kappa_1 \cdot \sqrt{\frac{\log(T-t+1)}{T-t+1}} \), then we set \( \hat{M}_{c,a_{j_t}} = l_{j_t} \).
4: else if \( \hat{q}_{j_t,t}^* \leq 3\kappa_1 \cdot \sqrt{\frac{\log(T-t+1)}{T-t+1}} \), then we set \( \hat{M}_{c,a_{j_t}} = u_{j_t} \).
5: else if \( 3\kappa_1 \cdot \sqrt{\frac{\log(T-t+1)}{T-t+1}} \leq \hat{q}_{j_t,t} \leq 1 - 3\kappa_1 \cdot \sqrt{\frac{\log(T-t+1)}{T-t+1}} \), then we set \( \hat{M}_{c,a_{j_t}} = F_{j_t}^{-1}(1 - \hat{q}_{j_t,t}^*) \).
6: Output: \( \hat{M}_{c,a_{j_t}} \)

Theorem 1. Suppose that the estimator \( \hat{M} \) is given in Algorithm 2. Then, as long as \( s = T - t + 1 \geq s_0 \) for a constant \( s_0 \geq 0 \) that solely depends on \( \{p_j\}_{j=1}^n \), it holds that
\[
\sup_{c \geq 0} \text{Myopic}_t(\pi, c) \leq \frac{\kappa_2 \cdot \log s}{s}
\]
for a constant \( \kappa_2 > 0 \).

Proof of Theorem 1. We have that
\[
V_{c}^{\text{Semi}}(\bar{d}_t) = \sum_{j=1}^{n} \hat{d}_{j,t+1} \cdot \int_{q=1-\hat{q}_{j,t}^*}^{1} F_{j}^{-1}(q) dq + \int_{q=1-\hat{q}_{j,t}^*}^{1} F_{j}^{-1}(q) dq.
\tag{16}
\]
where we denote by \( j_t \) the type of query \( t \) in the instance \( I_t \). Now we plug (16) into the formulation (8). We get
\[
\text{Myopic}_t(\pi, c) = \mathbb{E}_{\hat{M}_{c,a_{j_t}}} \left[ \sum_{j=1}^{n} \hat{d}_{j,t+1} \cdot \int_{q=1-\hat{q}_{j,t}^*}^{1} F_{j}^{-1}(q) dq + \int_{q=1-\hat{q}_{j,t}^*}^{1} F_{j}^{-1}(q) dq \right. \\
- \left. \int_{q=1-\hat{q}_{j,t}^*}^{1} F_{j}^{-1}(q) dq - \mathbb{E}_{r \sim F_{j_t}} \left[ V_{c}^{\text{Semi}}(\bar{d}_{j_t}^* r_{j_t^*}(t+1)) \right] \right]
\tag{17}
\]
where we denote by $\pi$ our online policy Algorithm 1 with the estimator given in Algorithm 2. Then, $q_{j,t}^\pi$ denotes the ex-ante probability that query $t$ will be served by the online policy $\pi$. With these notations, (17) can be re-written as

$$
\text{Myopic}_t(\pi, c) = E_{j_t,t+1} \left[ \sum_{j=1}^{n} \tilde{d}_{j,t+1} \cdot \int_{q=1-q_{j,t}^*}^{1} F_j^{-1}(q) dq + \int_{q=q_{j,t}^*}^{1} F_j^{-1}(q) dq - q_{j,t}^\pi \cdot \tilde{V}_{c-a_{j,t}}^{\text{Semi}}(\tilde{d}_{t+1}) - (1-q_{j,t}^\pi) \cdot \tilde{V}_{c-a_{j,t}}^{\text{Semi}}(\tilde{d}_{t+1}) \right]
$$

We construct feasible solution to $\tilde{q}_{j,t}^*$ for all the three cases. We identify a “good” event that $\tilde{q}_{j,t}^*$ is close to $q_{j,t}^*$ for each $j \in [n]$. Following Lemma 3, we know that

$$
|q_{j,t}^* - \tilde{q}_{j,t}^*| \leq \kappa_1 \cdot \max_{j'} \{p_{j'} - \tilde{d}_{j',t+1}/(s-1)\}.
$$

(18)

for a constant $\kappa_1 > 0$. We note that for each $j \in [n]$, $\tilde{d}_{j,t+1}$ is a binomial distribution with mean $p_j \cdot (s-1)$. Then, from Hoeffding’s inequality (Lemma 14), we have

$$
P(|\tilde{d}_{j,t+1} - p_j(s-1)| \leq \sqrt{4(s-1) \log(s-1)}) \geq 1 - 2 \exp(-2 \log(s-1)) \geq 1 - \frac{2}{(s-1)^2} \geq 1 - \frac{1}{n(s-1)}
$$

as long as $s \geq s_0$ for a constant $s_0$. We denote by the event

$$
\mathcal{G} = \{ |\tilde{d}_{j,t+1} - p_j(s-1)| \leq \sqrt{4(s-1) \log(s-1)}, \forall j \in [n]\},
$$

which is the “good” event that $\tilde{d}_{j,t+1}$ is close to its mean. From union bound, we know that

$$
P(\mathcal{G}) \geq 1 - \sum_{j=1}^{n} (1 - P(|\tilde{d}_{j,t+1} - p_j(s-1)| \leq \sqrt{4(s-1) \log(s-1)})) \geq 1 - \frac{1}{s-1}.
$$

We denote by

$$
\text{Myopic}_t(\pi, c, \mathcal{G}) = P(\mathcal{G}) \cdot E_{j_t,t+1} \left[ \sum_{j=1}^{n} \tilde{d}_{j,t+1} \cdot \int_{q=1-q_{j,t}^*}^{1} F_j^{-1}(q) dq + \int_{q=q_{j,t}^*}^{1} F_j^{-1}(q) dq - q_{j,t}^\pi \cdot \tilde{V}_{c-a_{j,t}}^{\text{Semi}}(\tilde{d}_{t+1}) - (1-q_{j,t}^\pi) \cdot \tilde{V}_{c-a_{j,t}}^{\text{Semi}}(\tilde{d}_{t+1}) \right | \mathcal{G}
$$

and

$$
\text{Myopic}_t(\pi, c, \mathcal{G}^c) = P(\mathcal{G}^c) \cdot E_{j_t,t+1} \left[ \sum_{j=1}^{n} \tilde{d}_{j,t+1} \cdot \int_{q=1-q_{j,t}^*}^{1} F_j^{-1}(q) dq + \int_{q=q_{j,t}^*}^{1} F_j^{-1}(q) dq - q_{j,t}^\pi \cdot \tilde{V}_{c-a_{j,t}}^{\text{Semi}}(\tilde{d}_{t+1}) - (1-q_{j,t}^\pi) \cdot \tilde{V}_{c-a_{j,t}}^{\text{Semi}}(\tilde{d}_{t+1}) \right | \mathcal{G}^c
$$

where $\mathcal{G}^c$ is the complement of $\mathcal{G}$. It is clear that

$$
\text{Myopic}_t(\pi, c) = \text{Myopic}_t(\pi, c, \mathcal{G}) + \text{Myopic}_t(\pi, c, \mathcal{G}^c)
$$
We have a direct upper bound
\[
\text{Myopic}_t(\pi, c, \mathcal{G}) \leq \frac{\max_{j \in [n]} \{ u_j \}}{s - 1}.
\]

In what follows, we condition on the event \( \mathcal{G} \) happens, and we bound \( \text{Myopic}_t(\pi, c, \mathcal{G}) \).

**Case (i) when \( q_{jt,t}^* \geq 1 - 3\kappa_1 \cdot \sqrt{\frac{\log s}{s}} \):** in this case, we know that
\[
c \geq p_{jt} \cdot s \cdot \hat{q}_{jt,t}^* \cdot a_{jt} \geq \frac{p_{jt} \cdot s \cdot a_{jt}}{2} \geq a_{jt}
\]
when \( s \geq s_0 \) for a constant \( s_0 \geq 0 \). Therefore, we always have enough remaining capacity to serve query \( t \) with type \( j_t \).

Since we have \( q_{jt,t}^* = 1 \), we only need to construct feasible solution to \( \tilde{V}_{c-a_{jt}}(\hat{d}_{t+1}) \). From the feasibility of \( \{ \hat{q}_{jt,t}^* \} \), we know that
\[
\sum_{j=1}^n \tilde{d}_{j,t+1} \cdot a_{jt,i} \cdot \hat{q}_{jt,t}^* + a_{jt,i} \cdot \tilde{q}_{jt,t}^* \leq c_i, \ \forall i \in [m].
\]

Note that conditioning on the event \( \mathcal{G} \), following (18), we have
\[
|\hat{q}_{jt,t}^* - \tilde{q}_{jt,t}^*| \leq 2\kappa_1 \cdot \sqrt{\frac{\log(s-1)}{s-1}}
\]
which implies \( \hat{q}_{jt,t}^* \geq 1 - 5\kappa_1 \cdot \sqrt{\frac{\log s}{s}} \geq \frac{1}{2} \) when \( s \geq s_0 \) for a constant \( s_0 > 0 \). We construct the following solution \( \{ \tilde{q}_{jt,t}^* \} \) for \( \tilde{V}_{c-a_{jt}}(\hat{d}_{t+1}) \) satisfying
\[
\tilde{q}_{jt,t}^* = \hat{q}_{jt,t}^*, \forall j \neq j_t \text{ and } \tilde{q}_{jt,t}^* = \hat{q}_{jt,t}^* + \frac{\tilde{q}_{jt,t}^* - 1}{\tilde{d}_{jt,t+1}}.
\]

Since \( \sum_{j=1}^n \tilde{d}_{j,t+1} \cdot a_{jt,i} \cdot \tilde{q}_{jt,t}^* \leq c_i - a_{jt,i} \) for each \( i \in [m] \), we know that \( \{ \tilde{q}_{jt,t}^* \} \) is a feasible solution to \( \tilde{V}_{c-a_{jt}}(I_{t+1}) \), where \( \tilde{q}_{jt,t}^* \geq 0 \) follows from \( \hat{q}_{jt,t}^* \geq \frac{1}{2} \) and \( \tilde{d}_{jt,t+1} \geq 1 \) conditioning on the event \( \mathcal{G} \).

Therefore, we have that
\[
\text{Myopic}_t(\pi, c, \mathcal{G}) \leq \mathbb{E}_{I_{t+1}} \left[ \sum_{j=1}^n \tilde{d}_{j,t+1} \cdot \int_{q=1-\tilde{q}_{jt,t}^*}^1 F_{jt}^{-1}(q) dq + \int_{q=1-\tilde{q}_{jt,t}^*}^{1-\hat{q}_{jt,t}^*} F_{jt}^{-1}(q) dq \right.
\]
\[
- \sum_{j=1}^n \tilde{d}_{j,t+1} \cdot \int_{q=1-\hat{q}_{jt,t}^*}^1 F_{jt}^{-1}(q) dq \mid \mathcal{G} \right]
\]
\[
= \mathbb{E}_{I_{t+1}} \left[ \tilde{d}_{jt,t+1} \cdot \int_{q=1-\tilde{q}_{jt,t}^*}^{1-\hat{q}_{jt,t}^*} F_{jt}^{-1}(q) dq + \int_{q=1-\tilde{q}_{jt,t}^*}^{1-\hat{q}_{jt,t}^*} F_{jt}^{-1}(q) dq \mid \mathcal{G} \right]
\]
\[
(22)
\]

We make the following claim.

**Claim 1.** For any \( q_1, q_2 \in [0,1] \), it holds that
\[
\int_{q=q_1}^{q_2} F_{jt}^{-1}(q) dq \leq F_{jt}^{-1}(q_1) \cdot (q_2 - q_1) + \frac{(q_2 - q_1)^2}{\alpha}
\]
for any \( j \in [n] \), where \( \alpha \) is the lower bound of the density function \( f(\cdot | a_j) \) specified in Assumption 1.
Since \( \sum_{q=1}^{q^*} \) Applying Claim 1, we have
\[
\int_{q=1}^{1-q^*} F_{jt}(q) dq \leq F_{jt}^{-1}(1-q^*) \cdot \frac{1-q^*}{d_{jt,t+1}} + \frac{(1-q^*)^2}{\alpha \cdot d_{jt,t+1}^2}
\] (23)
and
\[
\int_{q=1}^{1-q^*} F_{jt}(q) dq \leq F_{jt}^{-1}(1-q^*) \cdot (q^* - 1) + \frac{(1-q^*)^2}{\alpha}
\] (24)

Plugging (23) and (24) into (22), we have
\[
\text{Myopic}_t(\pi, c, G) \leq \mathbb{E}_{j,t+1} \left[ \left( \frac{(1-q^*)^2}{\alpha} + \frac{(1-q^*)^2}{\alpha \cdot d_{jt,t+1}} \right) | G \right] \cdot P(G)
\]
\[
\leq 2 \mathbb{E}_{j,t+1} \left[ \frac{(1-q^*)^2}{\alpha} + \frac{(q^* - \hat{q}_{jt,t})^2}{\alpha \cdot d_{jt,t+1}} \right] | G \right] \cdot P(G)
\]
\[
\leq 2 \log s + \frac{2}{\alpha} \mathbb{E}_{j,t+1}[(q^* - \hat{q}_{jt,t})^2] + \frac{2}{\alpha} \cdot \mathbb{E}_{j,t+1} \left[ \frac{1}{d_{jt,t+1}} \right] | G \right]
\]
(25)

**Case (ii) when \( \hat{q}_{jt,t}^* \leq 3\kappa_1 \cdot \sqrt{\frac{\log s}{s}} \):** since we have \( q_{jt,t}^* = 0 \), we only need to construct feasible solution to \( \tilde{V}_c^\text{semi} (\tilde{d}_{t+1}) \). Note that conditioning on the event \( G \), following (18), we have
\[
|q^* - \hat{q}^*| \leq 2\kappa_1 \cdot \sqrt{\frac{\log (s-1)}{s}}
\]
which implies \( \hat{q}_{jt,t}^* \leq 5\kappa_1 \cdot \sqrt{\frac{\log s}{s}} \leq \frac{1}{2} \) as long as \( s \geq s_0 \) for a constant \( s_0 \geq 0 \). From the feasibility of \( \{\hat{q}_{jt,t}^*\} \) demonstrated in (20), we construct the following solution \( \{\hat{q}_{jt,t}''\} \) for \( \tilde{V}_c^\text{semi} (\tilde{d}_{t+1}) \) satisfying
\[
\hat{q}_{jt,t}'' = \hat{q}_{jt,t}^*, \forall j \neq j_t \quad \text{and} \quad \hat{q}_{jt,t}'' = \hat{q}_{jt,t}^* \cdot \frac{d_{jt,t+1} + 1}{d_{jt,t+1}}
\]
(26)

Since \( \sum_{j=1}^{n} d_{jt+1} \cdot a_{jt} \cdot \hat{q}''_{jt,t} \leq c_i \) for each \( i \in [m] \), we know that \( \{\hat{q}_{jt,t}''\} \) is a feasible solution to \( \tilde{V}_c^\text{semi} (\tilde{d}_{t+1}) \), where \( \hat{q}_{jt,t}'' \leq 1 \) follows from \( \hat{q}_{jt,t}^* \leq \frac{1}{2} \) and \( d_{jt,t+1} \geq 1 \) conditioning on the event \( G \).

Therefore, we have that
\[
\text{Myopic}_t(\pi, c, G) \leq \mathbb{E}_{j,t+1} \left[ \sum_{j=1}^{n} d_{jt,t+1} \cdot \int_{q=1}^{1-q_{jt,t}^*} F_{jt}(q) dq + \int_{q=1}^{1-q_{jt,t}^*} F_{jt}^{-1}(q) dq \right] + \mathbb{E}_{j,t+1} \left[ \frac{1}{d_{jt,t+1}} \left( \int_{q=1}^{1-q_{jt,t}^*} F_{jt}^{-1}(q) dq \right) \right]
\]
(27)

Applying Claim 1, we have that
\[
\int_{q=1}^{1-q_{jt,t}^*} F_{jt}^{-1}(q) dq \leq -F_{jt}^{-1}(1-q_{jt,t}^*) \cdot \frac{\hat{q}_{jt,t}^*}{d_{jt,t+1}} + \frac{(\hat{q}_{jt,t}^*)^2}{\alpha \cdot d_{jt,t+1}^2}
\]
(28)
and
\[
\int_{q = 1 - \tilde{q}_{j,t}^*}^{1 - \mu_{j,t}} F_{j,t}^{-1}(q)dq \leq F_{j,t}^{-1}(1 - \tilde{q}_{j,t}^*) \cdot \tilde{q}_{j,t}^* + \frac{(\tilde{q}_{j,t}^*)^2}{\alpha} \tag{29}
\]

Plugging (28) and (29) into (27), we get
\[
\text{Myopic}(\pi, c, G) \leq \mathbb{E}_{j_t, t+1} \left[ \frac{(\tilde{q}_{j,t}^*)^2}{\alpha} + \frac{(\tilde{q}_{j,t}^*)^2}{\alpha \cdot \tilde{d}_{j,t+1}} \mid G \right] \cdot P(G)
\]
\[
\leq 2 \mathbb{E}_{j_t, t+1} \left[ \frac{(\tilde{q}_{j,t}^*)^2}{\alpha} + \frac{(\tilde{q}_{j,t}^* - \tilde{q}_{j,t}^*)^2}{\alpha} + \frac{1}{\alpha \cdot \tilde{d}_{j,t+1}} \mid G \right] \cdot P(G) \tag{30}
\]
\[
\leq \frac{2 \log s}{\alpha \cdot s} + \frac{2}{\alpha} \cdot \mathbb{E}_{j_t, t+1}[((\tilde{q}_{j,t}^* - \tilde{q}_{j,t}^*)^2] + \frac{2}{\alpha} \cdot \mathbb{E}_{j_t, t+1} \left[ \frac{1}{\tilde{d}_{j,t+1}} \mid G \right].
\]

**Case (iii) when** $3\kappa_1 \cdot \sqrt{\frac{\log s}{s}} \leq \tilde{q}_{j,t}^* \leq 1 - 3\kappa_1 \cdot \sqrt{\frac{\log s}{s}}$; in this case, we know that
\[
c \geq p_{j_t} \cdot s \cdot \tilde{q}_{j,t}^* \cdot \alpha_{j_t} \geq p_{j_t} \cdot \sqrt{s \log s} \cdot \alpha_{j_t} \geq \alpha_{j_t}
\]
when $s \geq s_0$ for a constant $s_0 \geq 0$ that depends solely on $\{p_j\}_{j=1}^n$. Therefore, we always have enough remaining capacity to serve query $t$ with type $j_t$.

Note that conditioning on the event $G$, following (18), we have
\[
|\tilde{q}_{j,t}^* - \tilde{q}_{j,t}^*| \leq 2\kappa_1 \cdot \sqrt{\frac{\log(s-1)}{s-1}}
\]
which implies $\kappa_1 \cdot \sqrt{\frac{\log s}{s}} \leq \tilde{q}_{j,t}^* \leq 1 - \kappa_1 \cdot \sqrt{\frac{\log s}{s}}$.

We construct feasible solution $\{\tilde{q}_{j,t}''\}$ for $V_{c - \alpha_{j_t}}^{Semi}(\tilde{d}_{t+1})$ following the definition in (21). Then, $\tilde{q}_{j,t}' \geq 0$ follows from the fact that
\[
\tilde{q}_{j,t}^* \geq \kappa_1 \cdot \sqrt{\frac{\log s}{s}} \geq \frac{2}{p_{j_t} \cdot (s-1)} \geq \frac{1}{\tilde{d}_{j,t+1}}
\]
-conditioning on the event $G$, as long as $s \geq s_0$ for a constant $s_0 > 0$.

We construct feasible solution $\{\tilde{q}_{j,t}''\}$ for $V_{c}^{Semi}(\tilde{d}_{t+1})$ following the definition in (26). Then, $\tilde{q}_{j,t}'' \leq 1$ follows from
\[
\tilde{q}_{j,t}'' \leq \tilde{q}_{j,t}^* + \frac{1}{\tilde{d}_{j,t+1}} \leq 1 - \kappa_1 \cdot \sqrt{\frac{\log s}{s}} + \frac{3}{2p_{j_t}(s-1)} \leq 1
\]
-conditioning on the event $G$, as long as $s \geq s_0$ for a constant $s_0 > 0$. 

Therefore, we have that
\[
\text{Myopic}_t(\pi, c, G) \leq \mathbb{E}_{j_t, t+1} \left[ \sum_{j=1}^{n} \tilde{d}_{j, t+1} \cdot \int_{q=1}^{1} F_{j_t}^{-1}(q) dq + \int_{q=1}^{1-\tilde{q}_{j_t}^*} F_{j_t}^{-1}(q) dq \right. \\
- q_{j_t}^* \cdot \sum_{j=1}^{n} \tilde{d}_{j, t+1} \cdot \int_{q=1}^{1} nF_{j_t}^{-1}(q) dq - (1-q_{j_t}^* \cdot \sum_{j=1}^{n} \tilde{d}_{j, t+1} \cdot \int_{q=1}^{1-\tilde{q}_{j_t}^*} F_{j_t}^{-1}(q) dq \\
\left. \mathbb{E}_{j_t, t+1} \left[ \int_{q=1}^{1-\tilde{q}_{j_t}^*} F_{j_t}^{-1}(q) dq \mid G \right] \cdot P(G) \right. \\
+ \int_{q=1}^{1-\tilde{q}_{j_t}^*} F_{j_t}^{-1}(q) dq \mid G \right] \cdot P(G).
\] (31)

Applying Claim 1, we get (23), (28), and
\[
\int_{q=1}^{1-\tilde{q}_{j_t}^*} F_{j_t}^{-1}(q) dq = F_{j_t}^{-1}(1-\tilde{q}_{j_t}^*) \cdot (\tilde{q}_{j_t}^* - q_{j_t}^*) + \frac{(\tilde{q}_{j_t}^* - q_{j_t}^*)^2}{\alpha}.
\] (32)

Therefore, we have
\[
\text{Myopic}_t(\pi, c, G) \leq \mathbb{E}_{j_t, t+1} \left[ q_{j_t}^* \cdot \frac{(1-\tilde{q}_{j_t}^*)^2}{\alpha \cdot \tilde{d}_{j, t+1}} + (1-q_{j_t}^*) \cdot \frac{(\tilde{q}_{j_t}^*)^2}{\alpha \cdot \tilde{d}_{j, t+1}} + \frac{(\tilde{q}_{j_t}^* - q_{j_t}^*)^2}{\alpha} \mid G \right] \cdot P(G)
\leq \mathbb{E}_{j_t, t+1} \left[ \frac{1}{\alpha \cdot \tilde{d}_{j, t+1}} + \frac{(\tilde{q}_{j_t}^* - q_{j_t}^*)^2}{\alpha} \mid G \right] \cdot P(G).
\] (33)

From (25), (30) and (33), for all cases, it holds that
\[
\text{Myopic}_t(\pi, c, G) \leq \frac{2 \log s}{\alpha \cdot s} + \frac{2}{\alpha} \cdot \mathbb{E}_{j_t, t+1} [(\tilde{q}_{j_t}^* - q_{j_t}^*)^2] + \frac{2}{\alpha} \cdot \mathbb{E}_{j_t, t+1} \left[ \frac{1}{\tilde{d}_{j, t+1}} | G \right] \leq \frac{2 \log s}{\alpha \cdot s} + \frac{2}{\alpha} \cdot \mathbb{E}_{j_t, t+1} [(\tilde{q}_{j_t}^* - q_{j_t}^*)^2] + \frac{4}{\alpha (s-1)} \cdot \sum_{j=1}^{n} \frac{1}{p_j}.
\]

From Lemma 3, we know that there exists a constant \( \kappa_2 \) such that
\[
\mathbb{E}_{j_t, t+1} [(\tilde{q}_{j_t}^* - q_{j_t}^*)^2] \leq \kappa_2 \cdot \sum_{j=1}^{n} (p_j - \tilde{d}_{j, t}/s)^2 \leq \frac{\kappa_2}{s}
\]

Therefore, we know that
\[
\text{Myopic}_t(\pi, c, G) \leq \frac{2 \log s}{\alpha \cdot s} + \frac{2 \kappa_2}{\alpha \cdot s} + \frac{4}{\alpha (s-1)} \cdot \sum_{j=1}^{n} \frac{1}{p_j}.
\]

Together with (19), our proof is completed.

Following Lemma 1 and Theorem 1, we can see that the regret of the policy \( \pi \) given in Algorithm 2 can be upper bounded by
\[
\text{Regret}(\pi) \leq \sum_{s=1}^{T} \frac{\kappa_2 \cdot \log s}{s} \leq \kappa_2 \cdot \log^2 T
\]
which is our final regret bound.
4. Policy with Logarithmic Regret under Second-order Growth

In this section, we derive an improved logarithmic regret bound for (6) under the following stronger assumption.

**Assumption 2.** We assume that for each \( j \in [n] \), conditional on \( \tilde{a}_t \) being realized as any \( a_j \in A \), the reward distribution of \( \tilde{r} \) is supported on the interval \([0, u_j]\) with a density function \( f(\cdot | a_j) \), where \( u_j \geq 0 \), and it satisfies \( \beta \geq f(r | a_j) \geq \alpha \), for a constant \( \beta \geq \alpha > 0 \), for any \( r \in [0, u_j] \).

Note that Assumption 2 is stronger than Assumption 1 in that we require \( l_j = 0 \) for each \( l_j \) in Assumption 1. Moreover, in Assumption 2, we require not only the lower bound \( \alpha \) on the density, but also an upper bound \( \beta \) on the density \( f(\cdot | a_j) \) for each \( j \in [n] \). Assumption 2 implies the following less restrictive assumption, which will be used in our proofs.

**Assumption 3.** There exists a compact convex set \( \Omega \subset \mathbb{R}^m \geq 0 \) such that for any \( t \in [T] \), any \( c \) and any problem instance \( I_t \), the relaxed offline optimum \( \bar{V}_{t,c}(I) \) (10) possesses one optimal dual solution \( \tilde{\mu} \) satisfying \( \tilde{\mu} \in \Omega \). Moreover, there exists two positive constants \( \underline{\alpha}, \bar{\alpha} \) such that for any \( \mu', \mu'' \in \Omega \), it holds that

\[
\underline{\alpha} \cdot \mathbb{E}_a \left[ (\tilde{a}^\top \mu' - \tilde{a}^\top \mu'')^2 \right] \leq \mathbb{E}_a \left[ (F(\tilde{a}^\top \mu' | \tilde{a}) - F(\tilde{a}^\top \mu'' | \tilde{a})) \cdot (\tilde{a}^\top \mu' - \tilde{a}^\top \mu'') \right] \leq \bar{\alpha} \cdot \mathbb{E}_a \left[ (\tilde{a}^\top \mu' - \tilde{a}^\top \mu'')^2 \right]
\]

(34)

where \( \tilde{a} \) is a random variable denoting the size of one query and is realized as \( a_j \) with probability \( p_j \), for \( j \in [n] \).

We now provide an example to illustrate under which condition Assumption 3 holds or not.

**Example 1.** Consider a special case of our model where there is a single resource, i.e., \( m = 1 \), and two types of queries, i.e., \( n = 2 \). For each type \( j = 1 \) or \( 2 \), the size \( a_j = 1 \) and the reward distribution is a uniform distribution over the interval \([l_j, u_j]\) with \( u_1 \leq u_2 \). Note that Assumption 3 essentially requires that \( u_1 \geq l_2 \), i.e., the support for the reward distribution of each type of query overlaps with each other. In order to see this point, the set \( \Omega \) can be specified as \([l_1, u_2]\) and it is clear to see that (34) will be satisfied. In contrast, if \( u_1 < l_2 \), then we can see that (34) will be violated by setting \( \mu' = l_2 \) and \( \mu'' = u_1 \). However, the results derived in Section 3 will still apply and a \( O(\log^2 T) \) regret bound can be obtained.

The only reason why we make Assumption 3 is to establish a second-order growth condition of a dual function (we formalize in Lemma 7), which is a standard condition in the stochastic programming literature. Then, we apply results from the stochastic programming literature to derive our logarithmic regret bound (we formalize in Lemma 10). Note that the second-order growth condition has been assumed frequently in the previous literature (e.g. Bray (2022), Li and Ye (2021), Balseiro et al. (2021)). By deriving our results simply under Assumption 3, our contribution...
would be to get rid of a so-called “non-degeneracy” assumption, which concerns the “position” of the ex-ante relaxation $\bar{V}_{1,C}^{\text{Fld}}$ and is different from the second-order growth condition. In the following part, we briefly illustrate the “non-degeneracy” and we provide a thorough comparison between our Assumption 3 and the assumptions made in previous literature in Appendix A.

**Comparison with assumptions made in previous literature.** Notably, a common assumption made in the previous literature (Balseiro et al. 2021, Li and Ye 2021, Bray 2022) regarding logarithmic regret with continuous reward distribution is about the “non-degeneracy” of the ex-ante relaxation $\bar{V}_{1,C}^{\text{Fld}}$. The “non-degeneracy” assumption not only requires the optimal solution $\bar{V}_{1,C}^{\text{Fld}}$ to be unique, but also requires *strict complementary slackness* condition to be satisfied by $\bar{V}_{1,C}^{\text{Fld}}$, i.e., for any binding resource constraint in the optimal solution of $\bar{V}_{1,C}^{\text{Fld}}$, the corresponding optimal dual variable must be strictly positive. In this way, the optimal basis of $\bar{V}_{1,C}^{\text{Fld}}$ will remain unchanged if $C$ is perturbed by a certain amount, which drives all the analysis in Balseiro et al. (2021), Li and Ye (2021), Bray (2022). In contrast, our Assumption 3 simply requires the support of the reward distribution of each type to “overlap” with each other, as shown in Example 1, and our goal of Assumption 3 is to establish the standard second-order growth condition of the dual function. We require nothing over the “position” of the ex-ante relaxation $\bar{V}_{1,C}^{\text{Fld}}$, including unique optimal solution condition, strict complementary slackness condition, and so on.

We now provide an example where the strict complementary slackness condition of the ex-ante relaxation $\bar{V}_{1,C}^{\text{Fld}}$ is violated, while our Assumption 3 is still satisfied.

**Example 2.** Consider a model with 3 resources, each with an initial capacity of $T \cdot \frac{2\epsilon(1+\epsilon)}{1+5\epsilon}$. There are 3 types of queries, denoted by $j = 1, 2, 3$. The size of each type of query is $a_1 = (0,1,1)$, $a_2 = (1,0,1)$ and $a_3 = (1,1,0)$. The reward distribution of each type of query is a uniform distribution over $[0,1]$ and the arrival probability is $p_1 = \frac{2\epsilon}{1+5\epsilon}$, $p_2 = \frac{1+\epsilon}{1+5\epsilon}$ and $p_3 = \frac{2\epsilon}{1+5\epsilon}$. Clearly, our Assumption 2 is satisfied (thus Assumption 3 satisfied). On the other hand, following (13), the ex-ante relaxation $\bar{V}_{1,C}^{\text{Fld}} / T$ can be formulated as:

$$
\max \sum_{j=1}^{3} p_j \cdot q_j \cdot \left(1 - \frac{q_j}{2}\right) \\
\text{s.t.} \quad p_2q_2 + p_3q_3 \leq \frac{2\epsilon(1+\epsilon)}{1+5\epsilon} \\
p_1q_1 + p_3q_3 \leq \frac{2\epsilon(1+\epsilon)}{1+5\epsilon} \\
p_1q_1 + p_2q_2 \leq \frac{2\epsilon(1+\epsilon)}{1+5\epsilon} \\
q_1, q_2, q_3 \in [0,1]
$$

Denote by $\mu_i$ the dual variable for constraint for resource $i$. Then, it is clear to see that the primal dual pair $\mu_1^* = \mu_3^* = \frac{1-\epsilon}{2}$, $\mu_2^* = 0$ and $q_1^* = q_3^* = \frac{1+\epsilon}{2}$, $q_2^* = \epsilon$ is optimal to (35) by checking that the
saddle point condition is satisfied, for any $\epsilon > 0$. However, while the resource constraint is binding for every $i = 1, 2, 3$, the optimal dual variable $\mu_i = 0$, which shows that the strict complementary slackness condition is not satisfied for every $\epsilon > 0$. Therefore, the “non-degeneracy” assumption is violated for (35).

4.1. Decomposition of the Myopic Regret

We now proceed to bound Myopic$_i(\pi, \tilde{x}_i^\tau)$ in (8) under Assumption 1 and Assumption 3. We use $\tilde{V}^{\text{Off}}_{\tilde{c}^\tau_i}(I_t)$ (10) as the benchmark, which is a LP relaxation of $V^{\text{Off}}(I)$ (4) and thus provides more tractability. We denote by $\{\tilde{x}_i^\tau\}$ one optimal solution of $\tilde{V}^{\text{Off}}_{\tilde{c}^\tau_i}(I_t)$ (10), where $\tilde{x}_i^\tau \in [0, 1]$ for each $\tau = t, \ldots, T$. Then, a gap arises from the fact that the online decision for our policy $\pi$ must be binary as required by problem formulation, while the optimal solution $\tilde{x}_i^\tau$ in the relaxation (10) can be fractional. In order to deal with this gap, we introduce a “rounded” relaxed offline optimum as an intermediate. To be specific, we denote by

$$\tilde{x}_i^\text{round} = \begin{cases} 1, & \text{if } \tilde{r}_i \geq M_{\tilde{c}^\tau_i}a_i(I_{t+1}) = \tilde{V}^{\text{Off}}_{\tilde{c}^\tau_i}(I_{t+1}) - \tilde{V}^{\text{Off}}_{\tilde{c}^\tau_i-a_i}(I_{t+1}) \text{ and } \tilde{c}_i^\tau \geq \tilde{a}_i \\ 0, & \text{otherwise.} \end{cases}$$

(36)

as a rounding of $\tilde{x}_i^\tau$. Then, we have that

$$\tilde{V}^{\text{Off}}_{\tilde{c}^\tau_i}(I_t) = \tilde{x}_i^\text{round} \cdot (\tilde{r}_i + \tilde{V}^{\text{Off}}_{\tilde{c}^\tau_i-a_i}(I_{t+1})) + (1 - \tilde{x}_i^\text{round})\tilde{V}^{\text{Off}}_{\tilde{c}^\tau_i}(I_{t+1}) + G_{\tilde{c}^\tau_i}(I_t)$$

(37)

where $G_{\tilde{c}^\tau_i}(I_t)$ denotes the gap caused by rounding $\tilde{x}_i^\tau$ to be $\tilde{x}_i^\text{round}$, which can be formulated as follows

$$G_{\tilde{c}^\tau_i}(I_t) = \tilde{r}_i(\tilde{x}_i^\tau - \tilde{x}_i^\text{round}) + \tilde{V}^{\text{Off}}_{\tilde{c}^\tau_i-a_i}(I_{t+1}) - \left(\tilde{x}_i^\text{round}\tilde{V}^{\text{Off}}_{\tilde{c}^\tau_i-a_i}(I_{t+1}) + (1 - \tilde{x}_i^\text{round})\tilde{V}^{\text{Off}}_{\tilde{c}^\tau_i}(I_{t+1})\right).$$

(38)

Introducing the rounded gap $G_{\tilde{c}^\tau_i}(I_t)$ allows us to further decompose the regret of our policy into three terms, as formalized in the following lemma, where the proof is relegated to Appendix E.

**Lemma 5.** For any $t \in [T]$, it holds that

$$\text{Myopic}_i(\pi, \tilde{x}_i^\tau) \leq 2\tilde{r}_i \cdot \mathbb{E}_{a_i}[1(\tilde{c}_i^\tau \geq \tilde{a}_i) \cdot \text{Var}(M_{\tilde{c}_i^\tau}a_i(I_{t+1})) + \mathbb{E}_{I_{t+1}}[G_{\tilde{c}_i^\tau}(I_t)] + 2\tilde{r}_i \cdot \mathbb{E}_{a_i}[1(\tilde{c}_i^\tau \geq \tilde{a}_i) \cdot \left(\tilde{M}_{\tilde{c}_i^\tau}a_i - \mathbb{E}_{I_{t+1}}[M_{\tilde{c}_i^\tau}a_i(I_{t+1})]\right)^2]$$

(39)

where the variance $\text{Var}$ is taken over the problem instance $I_{t+1}$.

A key step in deriving Lemma 5 is to utilize the relationship

$$\tilde{V}^{\text{Off}}_{\tilde{c}_i}(I_t) = \tilde{r}_i \cdot \tilde{x}_i^\tau + \tilde{V}^{\text{Off}}_{\tilde{c}_i-a_i}(I_{t+1}).$$

(40)

We note that the above backward induction holds only on each sample path $I$, which is the reason why we use a sample-path based relaxed offline optimum $\tilde{V}^{\text{Off}}_{\tilde{c}_i}(I_t)$ (10) as the benchmark.
In contrast, if we use a benchmark that is independent of the sample path $I$, for example, the deterministic relaxation $\bar{V}^{\text{Fld}}$ (11) as considered in Li and Ye (2021) and Balseiro et al. (2021). Then, we only have the following induction:

$$V_{t,c}^{\text{Fld}} = \mathbb{E}_{(\tilde{r},\tilde{a})} [\tilde{r} \cdot x^*(\tilde{r},\tilde{a})] + \bar{V}_{t+1,c-\mathbb{E}_{(\tilde{r},\tilde{a})}[\tilde{a} \cdot x^*(\tilde{r},\tilde{a})]}^{\text{Fld}},$$

If we plug (41) into the proof of Lemma 5 to derive a similar decomposition of the myopic term $\text{Myopic}_t(\pi, \tilde{c}_t^\tau)$ that is now defined with respect to $\bar{V}_{t,c}^{\text{Fld}}$, then we would have an additional term $V_{t+1,c-\mathbb{E}_{(\tilde{r},\tilde{a})}[\tilde{a} \cdot x^*(\tilde{r},\tilde{a})]}^{\text{Fld}} - \mathbb{E}_{(\tilde{r},\tilde{a})}(\bar{V}_{t+1,c-\tilde{a} \cdot x^*(\tilde{r},\tilde{a})}]$ showing up in the decomposition. This additional term would require a “non-degeneracy” assumption (we further discuss in Appendix A) to deal with, as shown in Li and Ye (2021) and Balseiro et al. (2021). The above discussion reveals the benefits of considering a sample-path based benchmark $\bar{V}_{t,c}^{\text{Off}}(I_t)$, which enables us to get rid of the “non-degeneracy” assumption at the first hand. The same idea of using the relaxation $\bar{V}_{c}^{\text{Off}}(I_t)$ and relationship (40) has been developed in Vera and Banerjee (2021) to get rid of the “non-degeneracy” assumption when the reward for each type is deterministic. We generalize this idea to allow the reward for each type having a continuous distribution.

For the RHS of (39), we refer the first term as the variation gap, the second term as the rounding gap, and the third term as the estimator gap. In what follows, we proceed to bound the three terms separately. Our goal is to show that each gap can be bounded at the order of $O(\frac{1}{\sqrt{T-t}})$, and these bounds together will imply a $O(\log T)$ bound for Algorithm 1.

Our analysis relies on considering the dual problem of $\bar{V}_{c}^{\text{Off}}(I_{t+1})$ and $\bar{V}_{c}^{\text{Off}}(\tilde{a}_t(I_{t+1}))$ that define $M_{\tilde{c}_t, \tilde{a}_t}(I_{t+1})$ in (8), for each fixed $\tilde{c}_t^\tau$ and $\tilde{a}_t$ satisfying $\tilde{c}_t^\tau \geq \tilde{a}_t$. Now for any $c \geq 0$, we introduce a dual variable $\mu$ for the constraints of $\bar{V}_{c}^{\text{Off}}(I_{t+1})$ (10) and we denote by the function

$$L_{c,I_{t+1}}^{\text{Off}}(\mu) := \max_{\tilde{x}_t \in [0,1]^{\forall \tau = t+1,\ldots,T}} \left( \frac{c}{s-1} \right)^\top \mu + \frac{1}{s-1} \sum_{\tau = t+1}^{T} [\tilde{r}_\tau - \tilde{a}_\tau^\top \mu] \cdot \tilde{x}_\tau$$

(42)

as the dual function of $\bar{V}_{c}^{\text{Off}}(I_{t+1})$, scaled by $\frac{1}{s-1}$, with $s = T-t+1$. We now proceed to bound the variation gap, the rounding gap, and the estimator gap. In Section 4.2, we show that bounding the first two gaps can be reduced to bounding a so-called “dual convergence”, which concerns the variance of one optimal dual variable for minimizing the dual function $L_{c,I_{t+1}}^{\text{Off}}(\mu)$ (42). Then, we propose our $\hat{M}$-estimator and bound the “dual convergence” in Section 4.3 to complete our final bound over the myopic regret.
4.2. Reduction to Dual Convergence

In this section, we show how to reduce bounding each term in (39) to bounding the “dual convergence”, and we also propose our $\hat{M}$–estimator.

We first bound the variation gap $\text{Var}(M_{c,a_t}(I_{t+1}))$ for each fixed $c$ and $a_t$ satisfying $c \geq a_t$. We denote by

$$\mu_1 \in \arg\min_{\mu \in \Omega} L_{c,t+1}^{\text{Off}}(\mu) \quad \text{and} \quad \mu_2 \in \arg\min_{\mu \in \Omega} L_{c-a_t,t+1}^{\text{Off}}(\mu)$$

(43)

Note that $\mu_1$ (resp. $\mu_2$) is one optimal dual variable of $\bar{V}_{c}^{\text{Off}}(I_{t+1})$ (resp. $\bar{V}_{c-a_t}^{\text{Off}}(I_{t+1})$). Also, $\mu_1$ and $\mu_2$ are random variables, where the randomness comes from the randomness of the problem instance $I_{t+1}$. The goal of introducing $\mu_1$ and $\mu_2$ is to lower-bound and upper-bound $M_{c,a_t}(I_{t+1})$ in (8), as in the following lemma.

**Lemma 6.** For any problem instance $I_{t+1}$ and any $c \geq a_t$, it holds that

$$a_t^\top \mu_1 \leq M_{c,a_t}(I_{t+1}) = V_{c}^{\text{Off}}(I_{t+1}) - V_{c-a_t}^{\text{Off}}(I_{t+1}) \leq a_t^\top \mu_2.$$  

(44)

where $\mu_1$ and $\mu_2$ are defined in (43).

We proceed to bound $\text{Var}(M_{c,a_t}(I_{t+1}))$ with the help of Lemma 6. We note that the function $L_{c,t+1}^{\text{Off}}(\mu)$ can be regarded as a sample average approximation of the following stochastic optimization problem:

$$\min_{\mu \in \Omega} L_{c,t+1}^{\text{Fld}}(\mu) := \left(\frac{c}{s-1}\right)^\top \mu + \mathbb{E}_{(\tilde{r},\tilde{a}) \sim F}[\tilde{r} - \tilde{a}^\top \mu]^+$$

(45)

with $s-1 = T - t$ samples. Our analysis relies on showing the second-order growth condition of the “limiting” dual function $L_{c,t+1}^{\text{Fld}}$ defined in (45). Denote by $S$ the subspace spanned by the resource consumption vector $a$ of each query:

$$S = \{ \sum_{j \in [n]} \alpha_j \cdot a_j : \forall \alpha_j \in \mathbb{R} \}$$

(46)

As a result of Assumption 3, we have the following second-order growth condition of the limiting dual function $L_{c,t+1}^{\text{Fld}}$ for any $c \geq 0$.

**Lemma 7.** For any $c \geq 0$, we denote by $\mu^* \in \arg\min_{\mu \in \Omega} L_{c,t+1}^{\text{Fld}}(\mu)$ such that $\mu^* \in \Omega$ with $\Omega$ specified in Assumption 3. We also denote by $P_S$ the projection of any vector to the subspace $S$ (defined in (46)) spanned by the resource consumption vector $a$ of each query. Then for any $\mu \in \Omega$, it holds that

$$L_{c,t+1}^{\text{Fld}}(\mu) - L_{c,t+1}^{\text{Fld}}(\mu^*) \geq \frac{\alpha \beta}{2} \cdot \|P_S(\mu - \mu^*)\|^2_2$$

where $\beta$ denotes the smallest positive eigenvalue of $\mathbb{E}_{\tilde{a}}[\tilde{a} \cdot \tilde{a}^\top]$. \[ \]
The proof is relegated to Appendix E. Denote by
\[ \hat{\mu}_1 \in \arg \min_{\mu \in \Omega} L_{c,t+1}^{\text{Fld}}(\mu) \quad \text{and} \quad \hat{\mu}_2 \in \arg \min_{\mu \in \Omega} L_{c-a_t,t+1}^{\text{Fld}}(\mu). \] (47)

Following classical results from sample average approximation for stochastic programming (Shapiro 1993), we know that \( \hat{\mu}_1 \) (resp. \( \hat{\mu}_2 \)) converges to \( \tilde{\mu}_1 \) (resp. \( \tilde{\mu}_2 \)) in probability, as \( s - 1 = T - t \to \infty \).

Thus, we can use \( \hat{\mu}_1 \) (resp. \( \hat{\mu}_2 \)) as an approximation of \( E[\tilde{\mu}_1] \) (resp. \( E[\tilde{\mu}_2] \)) and obtain a bound over \( \text{Var}(\tilde{\mu}_1) \) and \( \text{Var}(\tilde{\mu}_2) \), which finally implies an upper bound of \( \text{Var}(M_{c,a_t}(I_{t+1})) \). We summarize the above arguments in the following lemma, which shows that we can reduce bounding the variation gap to bounding the “dual convergence” terms \( \text{E}_{I_{t+1}}[(a_i^\top \tilde{\mu}_1 - a_i^\top \hat{\mu}_1)^2] \) and \( \text{E}_{I_{t+1}}[(a_i^\top \tilde{\mu}_2 - a_i^\top \hat{\mu}_2)^2] \).

The proof is relegated to Appendix E.

**Lemma 8.** There exists a constant \( \kappa_1 > 0 \) such that for any \( c \geq 0 \) and any \( a_t \) satisfying \( c \geq a_t \), it holds that
\[ \text{Var}(M_{c,a_t}(I_{t+1})) \leq \frac{\kappa_1}{s-1} + \kappa_1 \cdot \text{E}_{I_{t+1}}[(a_i^\top \tilde{\mu}_1 - a_i^\top \hat{\mu}_1)^2] + \kappa_1 \cdot \text{E}_{I_{t+1}}[(a_i^\top \tilde{\mu}_2 - a_i^\top \hat{\mu}_2)^2] \]
with \( s = T - t + 1 \), where \( \tilde{\mu}_1 \) and \( \tilde{\mu}_2 \) are defined in (43), \( \hat{\mu}_1 \) and \( \hat{\mu}_2 \) are defined in (47), and \( \text{E}_{I_{t+1}}[\cdot] \) denotes taking expectation over \( I_{t+1} \) that decides the value of \( \tilde{\mu}_1 \) and \( \tilde{\mu}_2 \).

we then reduce bounding the rounding gap (second term in (39)) to bounding the “dual convergence”. Following previous notations, we denote by \( \{\tilde{x}^*_\tau\}_{\tau=1}^T \) an optimal solution to \( V^\text{Off}_c(I_t) \) (10). Note that \( \tilde{x}^*_\tau \in [0,1] \) can be fractional. We denote by \( \tilde{x}^\text{round}_i \) as the rounding of \( \tilde{x}^*_i \) as in (36). We bound the rounding gap \( \sum_{\tau=1}^T \text{E}[G^\tilde{e}_\tau(I_t)] \), where the formulation of \( G^\tilde{e}_\tau(I_t) \) is given in (38). Clearly, when \( \tilde{x}^*_i = \tilde{x}^\text{round}_i \), \( G^\tilde{e}_\tau(I_t) = 0 \). Thus, the rounding gap arises from the fact that \( \tilde{x}^*_i \in [0,1] \) can be different from \( \tilde{x}^\text{round}_i \in (0,1) \).

A key observation can be summarized as follows: i) if \( c \geq a_t \), then both \( \tilde{\mu}_1 \) and \( \tilde{\mu}_2 \) are well-defined in (43) and \( \tilde{x}^*_i \neq \tilde{x}^\text{round}_i \) happens if only \( \tilde{a}_i^\top \tilde{\mu}_1 \leq \tilde{r}_i \leq \tilde{a}_i^\top \tilde{\mu}_2 \); ii) if \( c \geq a_t \) does not hold and only \( \tilde{\mu}_1 \) is well-defined in (43), then \( \tilde{x}^*_i \neq \tilde{x}^\text{round}_i \) happens if only \( \tilde{a}_i^\top \tilde{\mu}_1 \leq \tilde{r}_i \). Therefore, we can again involve the dual variables \( \tilde{\mu}_1 \) and \( \tilde{\mu}_2 \) (if well-defined) in bounding the rounding gap \( G_c(I_t) \) for any \( c \geq 0 \), which reduces bounding the rounding gap into bounding the “dual convergence”. We formalize the above arguments in the following lemma, where the proof is relegated to Appendix E.

**Lemma 9.** For any \( c \), it holds that
\[ \text{E}_{I_t}[G_c(I_t)] \leq \frac{\kappa_2}{s-1} + \kappa_2 \cdot \text{E}_{a_t} \left[ \text{E}_{I_{t+1}}[(\tilde{a}_i^\top \tilde{\mu}_1 - \tilde{a}_i^\top \hat{\mu}_1)^2] \right] + \kappa_2 \cdot \text{E}_{a_t} \left[ \sum_{c \geq a_t} \text{E}_{I_{t+1}}[(\tilde{a}_i^\top \tilde{\mu}_2 - \tilde{a}_i^\top \hat{\mu}_2)^2] \right] \]
for a constant \( \kappa_2 \) that is independent of \( c \) and \( t \), where \( \tilde{\mu}_1 \) and \( \tilde{\mu}_2 \) (if \( c \geq a_t \) and well-defined) are defined in (43).
4.3. Bound on Dual Convergence and Policy

In this section, we first bound the “dual convergence” and then propose our $\hat{M}$—estimator. By further utilizing the second-order growth condition established in Lemma 7, we can show the following bound over the term $E[(a^T_t \tilde{\mu}_1 - a^T_t \hat{\mu}_1)^2]$. We only state the “dual convergence” result for $\hat{\mu}_1$ and $\tilde{\mu}_1$. Since the bound is independent of $c$, it is clear that the same bound also holds for $\hat{\mu}_2$ and $\tilde{\mu}_2$ whenever well-defined (i.e. $c \geq \tilde{a}_t$).

**Lemma 10.** For any $c \geq 0$, let $\hat{\mu}_1$ be defined in (47) and $\tilde{\mu}_1$ be defined in (43). Then, it holds that

$$E[(a^T_t \tilde{\mu}_1 - a^T_t \hat{\mu}_1)^2] \leq \frac{\delta_1}{s-1}$$

as long as $s-1 = T-t \geq t_0$, where $\delta_1$ and $t_0$ are constants that are independent of $s$ and $c$.

The key of the proof of Lemma 10 is to regard $\tilde{\mu}_1$ as the solution of the sample average approximation (definition in (43)) of the stochastic programming (45) with parameter $c$, whose optimal solution is $\hat{\mu}_1$ by definition (47). The second-order growth condition in Lemma 7 enables us to apply Theorem 2.1 in Shapiro (1993) showing that the gap between $\tilde{\mu}_1$ and $\hat{\mu}_1$ is at the order of $\sqrt{\frac{1}{s-1}}$ with high probability, as $s-1 = T-t \to \infty$. This result is also known as asymptotic normality of sample average solution in the stochastic programming literature. Note that in order to apply Theorem 2.1 in Shapiro (1993), there are some additional conditions need to be satisfied. We verified that all these conditions are satisfied by our problem in the proof of Lemma 10, which is relegated to Appendix E. However, the high probability result cannot be directly translated into a bound over the $L_2$ norm. In order to obtain the bound over the $L_2$ norm, we further utilize the method developed in Li and Ye (2021). As a result, a combination of the methods from Shapiro (1993) and Li and Ye (2021) enables us to bound the $L_2$ norm at the order of $\frac{1}{s-1}$, which improves the $\frac{\log \log(s-1)}{s-1}$ bound established in Theorem 1 in Li and Ye (2021).

We now present our $\hat{M}$—estimator to complete our algorithm and provide the final regret bound. From Lemma 6, we know that $M_{c,a_t}(I_{t+1})$ is close to $a^T_t \hat{\mu}_1$ and from the “dual convergence” established in Lemma 10, we know that $\hat{\mu}_1$ is a good “approximate” of $\tilde{\mu}_1$. Therefore, we use $a^T_t \hat{\mu}_1$ as our $\hat{M}$—estimator, which is formalized in Algorithm 1. We provide the following regret bound for the $\hat{M}$—estimator in Algorithm 3 and the proof is relegated to Appendix E.

**Theorem 2.** Suppose that the estimator $\hat{M}$ is given in Algorithm 3 and we denote by $\pi$ the policy given in Algorithm 1. Then, it holds that

$$E_{d^P_{c}} [\text{Myopic}_c(\pi, \hat{c}_t^P)] \leq \frac{\kappa_3}{T-t}$$

for a constant $\kappa_3 > 0$ that is independent of $t$. 
Algorithm 3 Algorithm achieving $O(\log T)$ Regret

1: Input: the remaining inventory $\tilde{c}_t^\pi$ and size $\tilde{a}_t$.

2: Obtain $\tilde{\mu}_1$ by solving

$$\min_{\mu \in \Omega} L_{\tilde{e}_t, t+1}^{\text{Fld}}(\mu) := \left( \frac{\tilde{c}_t^\pi}{T-t} \right)^T \mu + E_{(\tilde{r}, \tilde{a}) \sim F}[\tilde{r} - \tilde{a}^T \mu]^+$$

3: Output: $\hat{M}_{\tilde{e}_t, a_t} = a_t^T \tilde{\mu}_1$.

Therefore, by applying Lemma 1 and Theorem 2, we can see that the regret of the policy given in Algorithm 3 is upper bounded by $O(\log T)$.

Acknowledgments

The authors thank Rob Bray for motivations for removing the degeneracy assumption, and He Wang for bringing up the connection with price-based NRM.

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Appendix A: Detailed Comparison to Assumptions in Existing Literature

We conclude by making comparisons between our assumptions and the assumptions made in the existing literature, which can be summarized into the following two conditions:

(i). The non-degeneracy assumption over the ex-ante relaxation $\bar{V}^{\text{Fld}}$, which requires the optimal solution to $\bar{V}^{\text{Fld}}$ to be unique and strict complementary slackness condition being satisfied.

(ii). The second-order growth condition over the dual function $L^{\text{Fld}}$.

We first summarize the assumptions made in Li and Ye (2021) in the language of our paper as follows:

ASSUMPTION 4 (Assumptions made in Li and Ye (2021)). The following conditions have to be satisfied:

(i). $\tilde{r}_t$ and $\|\tilde{a}_t\|_2$ are always bounded for each $t \in [T]$.

(ii). $C$ scales linearly in $T$ and $C_i/T \in (d, \bar{d})$ with $\bar{d} \geq d > 0$, for each $i \in [m]$.

(iii). The matrix $E_{\tilde{a}}[\tilde{a}\tilde{a}^\top]$ is positive definite.

(iv). There exists a set $\Omega_{\mu}$ containing all possible optimal dual variable to the offline optimum $\bar{V}^{\text{Off}}$ (10) such that for any $\mu \in \Omega_{\mu}$ and any $c \in \Omega_c = [\bar{d} \cdot T, \bar{d} \cdot T]^m$, it holds that

$$\alpha \cdot |a^\top \mu - a^\top \mu^*(c)| \leq |F(a^\top \mu|a) - F(a^\top \mu^*(c)|a)| \leq \alpha \cdot |a^\top \mu - a^\top \mu^*(c)|$$

for any $a \in A$, where $\mu^*(c)$ is the optimal dual solution to the ex-ante relaxation $\bar{V}^{\text{Fld}}_{\mu,c}$ (11).

(v). For any $c \in \Omega$, the optimal dual solution $\mu^*(c)$ to the ex-ante relaxation $\bar{V}^{\text{Fld}}_{\mu,c}$ (11) satisfies the strict complementary slackness condition.

It is shown in Proposition 2 of Li and Ye (2021) that condition (iv) in Assumption 4 implies the second-order growth condition of the dual function, while condition (iii), (iv) and (v) all together imply the non-degeneracy assumption. We summarize the assumptions made in Bray (2022) as follows:

ASSUMPTION 5 (Assumptions made in Bray (2022)). The following conditions have to be satisfied:

(i). We have $f(r|a) \leq \bar{a}$ for any $a \in A$ and any $r$ in the support.

(ii). For any $a \in A$, $E_{(\tilde{r}, \tilde{a}) \sim F} [\tilde{r} | \tilde{a} = a] \leq \beta$ for a constant $\beta > 0$.

(iii). $\|a\|_2 \leq \tilde{d}$ for any $a \in A$.

(iv). The optimal dual solution to the Lagrangian problem $\min_{\mu \geq 0} L^{\text{Fld}}_{C:1}$ of the ex-ante relaxation $\bar{V}^{\text{Fld}}_{\mu,c}$ (11) is unique and the strictly complementary slackness condition is satisfied, when $c$ belongs to a neighborhood of $C$, which scales linearly in $T$.

(v). The Hessian matrix of $L^{\text{Fld}}_{C:1}(\mu^*)$ over $\mu^*$, where $\mu^* = \arg\min_{\mu \geq 0} L^{\text{Fld}}_{C:1}$ is full rank (equivalently, positive definite).

(vi). The Hessian matrix of $L^{\text{Fld}}_{C:1}(\mu)$ over $\mu$ is Lipschitz continuous when $\mu$ belongs to a neighborhood of $\mu^* = \arg\min_{\mu \geq 0} L^{\text{Fld}}_{C:1}$.

Note that condition (v) and (vi) in Assumption 5 together imply the second-order growth condition over the dual function, while the non-degeneracy assumption is stated in condition (iv) in Assumption 5. In particular, Balseiro et al. (2021) has summarized the assumptions in to the following two conditions:
Assumption 6 (Assumption 2 in Balseiro et al. (2021)). The following conditions have to be satisfied:

(i). The binding constraints for $V_{1,c}^{FLd}$ remains the same as the binding constraints for $V_{1,c}^{FLd}$, as long as $c$ belongs to a neighborhood of the initial capacity $C$, where we denote by $J$ the set of resource constraints that are binding.

(ii). There exists a constant $\kappa > 0$ such that

$$
V_{1,c}^{FLd} - V_{1,c}^{FLd} \geq \left( \nabla_c V_{1,c}^{FLd} \right)^\top \left( \frac{c}{T} - \frac{C}{T} \right) - \kappa \cdot \left( \frac{c_J}{T} - \frac{C_J}{T} \right)^2
$$

for all $c$ belonging to a neighborhood of the initial capacity $C$, where $c_J$ denotes $J$ part of the vector $c$ and $C_J$ denotes $J$ part of the vector $C$.

It has been shown that a sufficient condition to guarantee condition (i) in Assumption 6 is that strict complementary slackness condition is satisfied by $V_{1,c}^{FLd}$ (SC 8 in Balseiro et al. (2021)), and a sufficient condition to guarantee condition (ii) in Assumption 6 is that the second-order growth condition is satisfied by dual function $L^{FLd}$ (SC 7 in Balseiro et al. (2021)).

In Section 4, we need the second-order growth condition, but without the non-degeneracy condition, and we derive a $O(\log T)$ regret bound following our myopic regret framework. In constrast, in Section 3, we get rid of both the second-order growth condition and the non-degeneracy condition and consider our problem under the most general setting. Our main result is a $O(\log^2 T)$ regret bound. Both regret bounds are new in the literature.

Appendix B: Capturing the Price-based NRM Problem

In this appendix we explain how our regret results for NRM with accept/reject decisions (often called the quantity-based NRM problem) extend to the price-based NRM problem, and why our model (with a finite number of possible demand vectors but an infinite number of possible reward values) is quite natural for the pricing problem.

In the prototypical price-based NRM problem, a firm starts with initial resource vector $C \in \mathbb{R}^m$. The firm is selling a finite set of products, indexed $j = 1, \ldots, n$, with each product $j$ needing to consume resource vector $a_j \in \mathbb{R}^m$ in order to be sold. At each time $t = 1, \ldots, T$, the firm must first post a price $P^t_j$ for each product $j$ ($P^t_j$ can be $\infty$ if there are insufficient resources for product $j$ at time $t$). Afterward, customer $t$ arrives, wanting a random product $\tilde{h}_t$ drawn independently according to a known probability vector $(p_1, \ldots, p_n)$. Conditional on the customer wanting product $j$, their valuation $\tilde{v}_t$ is drawn independently from a known, $j$-specific distribution $H_j$ that is continuous over an interval $[l_j, u_j]$ with PDF $h_j$ satisfying $h_j(v) > 0$. Customer $t$ makes a purchase if and only if $\tilde{v}_t \geq P^t_{\tilde{h}_t}$, in which case the firm collects revenue $P^t_{\tilde{h}_t}$ and consumes resources $a_{\tilde{h}_t}$. The firm’s objective is to maximize the expected total revenue collected.

Definition 1. For all $j = 1, \ldots, n$, the virtual valuation corresponding to a valuation $v$ drawn from $H_j$ is $\phi_j(v) := v - \frac{1 - h_j(v)}{h_j(v)}$. The expected virtual surplus is defined to be

$$
E \left[ \max_{S \subseteq [T]}: \sum_{i \in S} a_{\tilde{h}_t} \leq C \phi_j(\tilde{v}_t) \right].
$$

For all $j$, let $F_j$ denote the CDF of the random variable $\phi_j(\tilde{v})$ when $\tilde{v}$ is drawn from $H_j$. 

The virtual valuation $\phi_j(v)$ is well-defined for all $v \in [l_j, u_j]$, by the assumption that $h_j(v) > 0$. We now make the standard regularity assumption that the virtual valuation function $\phi_j$ is monotonic, along with an assumption on $F_j$ having a lower-bounded PDF which will allow us to apply our results by considering accept/reject decisions on rewards drawn from the distribution $F_1, \ldots, F_n$.

**Assumption 7.** For all $j = 1, \ldots, n$, valuation distribution $H_j$ is such that $\phi_j$ is a non-decreasing function over $[l_j, u_j]$. Moreover, the virtual valuation distribution $F_j$ has a PDF $f_j$ satisfying $f_j(r) \geq \alpha$ for all values $r \geq 0$ lying in the support of $f_j$, where $\alpha > 0$ is a constant.

We note that even under the first part of Assumption 7 (regularity), virtual valuations can be negative. The second part of Assumption 7 imposes that the non-negative part of the support of virtual valuations for any $j = 1, \ldots, n$ has a PDF that is lower-bounded by $\alpha$.

**Theorem 3** (Myerson (1981), Chawla et al. (2010)). Under Assumption 7, the expected revenue of the optimal Bayesian incentive-compatible and individually-rational mechanism, which is an upper bound on the revenue of any online pricing policy, is equal to the expected virtual surplus. Moreover, consider the quantity-based NRM problem on rewards drawn from the virtual valuation distributions $F_1, \ldots, F_n$. Any online accept/reject policy for this problem can be converted into an online pricing policy for the price-based NRM problem, such that the expected revenue earned in the price-based NRM problem equals the expected reward collected in the quantity-based NRM problem.

Theorem 3 shows that regret guarantees for quantity-based NRM imply the same guarantees for price-based NRM. We note that “regret” in the price-based NRM setting is defined against the optimal mechanism benchmark, which is an upper bound on the revenue of any online pricing policy because these are special cases of sequential mechanisms.

When applying the accept/reject policy from the virtual valuation space on the original price-based NRM problem, for any $j = 1, \ldots, n$, one takes the minimum acceptable virtual valuation threshold $Q_j$ and maps it to a price $P_j \in [l_j, u_j]$ with the same probability of sale (i.e. $1 - F_j(Q_j) = 1 - H_j(P_j)$) and the same immediate expected reward/revenue (i.e. $\int_{Q_j}^{1} F_j^{-1}(q) dq = P_j(1 - H_j(P_j))$). Since virtual valuations can be negative, the accept/reject problem can face negative rewards, something not captured by our initial quantity-based NRM model. However, this is easily assuaged by splitting each “type” $j$ into two, one of which has a reward that is deterministically 0 (such rewards would never be accepted, so it is without loss of generality to convert negative virtual valuations into 0 rewards).

We now show that Assumption 7 is satisfied for many valuation distributions commonly used in the pricing literature, leading to the following corollary.

**Corollary 1.** If all valuation distributions satisfy Assumption 7, then logarithmic-level regret can be achieved for the price-based NRM problem. In particular, the following classes of valuation distributions satisfy Assumption 7:

1. Uniform valuations, i.e. $h_j(v) = \frac{1}{u_j - l_j}$ for all $v \in [l_j, u_j]$;
2. Truncated normal valuations, i.e. $h_j(v) = \alpha_j \cdot \exp(-v^2/\sigma_j)$ for all $v \in [l_j, u_j]$;
3. Truncated exponential valuations, i.e. $h_j(v) = \alpha_j \cdot \exp(-\lambda_j \cdot v)$ for all $v \in [l_j, u_j]$;
Proof.
1. Virtual valuations, defined by \( \phi_j(v) = 2v - u_j \) for all \( v \in [l_j, u_j] \), are uniformly distributed over \([2l_j - u_j, u_j]\), i.e. \( f_j(r) = \frac{1}{2(u_j - l_j)} \) for all \( r \in [2l_j - u_j, u_j] \);

2. Virtual valuations of the truncated normal valuations are known to be monotonic (Hartline 2013). We only prove the bounds for the pdf \( f_j(r) \). For a \( v_1, v_2 \in [l_j, u_j] \), we denote by

\[
r_1 = \phi_j(v_1) = v_1 - \frac{1 - H_j(v_1)}{h_j(v_1)} \quad \text{and} \quad r_2 = \phi_j(v_2) = v_2 - \frac{1 - H_j(v_2)}{h_j(v_2)}.
\]

Then, it holds that

\[
f_j(r_1) = \lim_{r_2 \to r_1} \frac{F_j(r_2) - F_j(r_1)}{r_2 - r_1} = \lim_{r_2 \to r_1} \frac{H_j(v_2) - H_j(v_1)}{\phi_j'(v_1) \cdot (v_2 - v_1)} \cdot h_j'(v_1) = f_j(v_1) \cdot h_j'(v_1) \frac{h_j(v_1)}{2}.
\]

By noting that \( h_j'(v_1) \leq 0 \), it is easy to see that

\[
\frac{h_j(v_1)}{2} \leq f_j(r_1)
\]

Clearly, \( \frac{h_j(v_1)}{2} \) is lower bounded by a positive constant on the interval \( v_1 \in [l_j, u_j] \), which implies that the pdf of the virtual valuation has a positive lower bound.

3. Virtual valuations of the truncated exponential valuations are known to be monotonic (Hartline 2013). We only prove the bounds for the pdf \( f_j(r) \). Following (48) and \( h_j'(v_1) \leq 0 \), we have

\[
\frac{h_j(v_1)}{2} \leq f_j(r_1).
\]

Therefore, \( f_j(r_1) \) is lower bounded by a positive constant for any \( r_1 \) by noting that \( h_j(v_1) \) is lower bounded by a positive constant for any \( v_1 \in [l_j, u_j] \).

\( \square \)

Appendix C: Useful Known Results

We proceed now to establish Lipschitz continuity of solutions of linear systems with respect to right-hand side perturbations.

**Lemma 11 (Theorem 2.2 of Mangasarian and Shiau (1987)).** Consider two linear systems

\[
\hat{A}x \leq b^1
\]

and

\[
\hat{A}x \leq b^2.
\]

For any solution \( x^1 \) that satisfies linear system (49), there exists a solution \( x^2 \) that satisfies linear system (50) such that

\[
\|x^1 - x^2\|_\infty \leq \mu \cdot \|b^1 - b^2\|_\infty
\]

where

\[
\mu = \sup_v \left\{ \left\| v^\top \hat{A} \right\|_1 \mid \left\| v \right\|_1 = 1, v \geq 0 \quad \text{Rows of } \hat{A} \text{ corresponding non-zero elements} \right\}.
\]

of \( v \) are linear independent.
We state the Hadamard’s inequality on matrix determinant in the following lemma.

**Lemma 12 (Hadamard (1893)).** For any matrix $A$, let $K$ denotes the number of columns in matrix $A$ and let $\{a_k\}_{k=1}^K$ denotes all the columns of matrix $A$. Then, the determinant of matrix $A$, denoted by $\det(A)$, satisfies the following inequality 

$$\det(A) \leq \prod_{k=1}^K \|a_k\|_2.$$ 

We state the well-known Bernstein’s inequality in the following lemma.

**Lemma 13 (Bernstein’s Inequality).** Let $X_1, \ldots, X_K$ be independent zero-mean random variables. Suppose that $|X_k| \leq M$ almost surely for all $k \in [K]$. Then, for all positive $\epsilon > 0$, it holds that 

$$P \left( \frac{1}{K} \sum_{i=1}^K X_k \geq \epsilon \right) \leq \exp \left( - \frac{1}{2} \frac{M^2 \epsilon^2}{\sum_{k=1}^K \text{Var}(X_k) + \frac{1}{4} K M \epsilon} \right).$$

We also state the well-known Hoeffdiding’s inequality in the following lemma.

**Lemma 14 (Hoeffding’s Inequality).** Let $X_1, \ldots, X_K$ be independent random variables such that $a_k \leq X_k \leq b_k$ almost surely, for each $k \in [K]$. Denote by $S_K = \frac{1}{K} \sum_{k=1}^K X_k$. Then, for any $\epsilon > 0$, it holds that 

$$P (|S_k - E[S_k] | \geq \epsilon) \leq 2 \exp(- \frac{2K^2 \epsilon^2}{\sum_{k=1}^K (b_k - a_k)^2}).$$

We then state the results from Huber (1967) under our notations.

**Assumption 8.** Suppose the following conditions hold:

(N-1). For each fixed $\mu$, the function $\psi((r,a), \mu)$ is separable.

(N-2). Denote $\lambda(\mu) = E_{(r,a)}[\psi((r,a), \mu)]$, then we have $\lambda(\mu^*) = 0$.

Denote 

$$u((r,a), \mu, d) = \sup_{\|\mu^* - \mu\|_2 \leq d} |\psi((r,a), \mu^*) - \psi((r,a), \mu)|$$

(N-3). There are strictly positive numbers $a, b, c_1, c_2, d_0$ such that

(i). $|\lambda(\mu)| \geq a \cdot \|\mu - \mu^*\|$ for $\|\mu - \mu^*\| \leq d_0$.

(ii). $E_{(r,a)}[u((r,a), \mu, d)] \leq b \cdot d$ for $\|\mu - \mu^*\| + d \leq d_0$.

(iii). $E_{(r,a)}[u((r,a), \mu, d)^2] \leq c_1 \cdot d$ for $\|\mu - \mu^*\| + d \leq d_0$.

(iv). $u((r,a), \mu, d) \leq c_2$ for any $(r,a)$.

Denote by 

$$Z_n(\mu^*, \mu'') = \frac{\sum_{j=1}^n [\psi((r_j, a_j), \mu^*) - \psi((r_j, a_j), \mu'') + \lambda(\mu') - \lambda(\mu'')]}{\sqrt{n} + n |\lambda(\mu')|}.$$ 

Then we have the following result from Huber (1967). Note that the original statement in Huber (1967) only concerns the convergence of $\sup_{\|\mu - \mu^*\| \leq d_0} Z_n(\mu, \mu^*)$ to 0 as $n \to \infty$. We now specify the constant terms in their bound and characterize the convergence rate, which will be helpful in our other proofs. The proof simply follows the proof in Huber (1967), except that we make specific the constant terms, and we include here for completeness.
**Lemma 15 (Lemma 3 in Huber (1967)).** The conditions in Assumption 8 imply that

\[
     \sup_{\|\mu - \mu^*\| \leq d_0} Z_n(\mu, \mu^*) \to 0
\]

in probability as \( n \to \infty \). Moreover, for any \( \epsilon > 0 \), it holds that

\[
    P \left( \sup_{\|\mu - \mu^*\| \leq d_0} Z_n(\mu, \mu^*) \geq 2\epsilon \right) \leq c \cdot e^{-n^{-\gamma'} + n^{-1}} \cdot \left( \frac{c}{b^2 q (1 - q)} + \frac{c}{9b^2 q^2 (1 - q)^2} \right) \cdot \left( \frac{\gamma' \log n}{\log(1 - q)} + 1 \right) \cdot (2M)^n
\]

as long as \( n \geq n_0 \), where \( n_0 \) satisfying \( n_0^{-\gamma'} = \frac{3b}{q} \) and \( \gamma' \in (\frac{1}{2}, 1) \) is an arbitrary number. Moreover, we set \( M \geq (3b)/(\epsilon a) \) and \( q = 1/M \).

**Proof of Lemma 15.** For the sake of simplicity, and without loss of generality, we choose the coordinate system such that \( \mu^* = 0 \). We also use \( x_j \) to denote \((r_j, a_j)\) for each \( j \in [n] \). The idea of the proof is to divide the cube \( \|\mu\| \leq d_0 \) into a slowly increasing number of smaller cubes and to bound \( Z_n(\mu, 0) \) in probability on each of those smaller cubes.

Put \( q = 1/M \), where \( M \geq 2 \) is an integer to be chosen later, and consider the concentric cubes

\[
    C_k = \{ \mu : \|\mu\| \leq (1 - q)^k \cdot d_0 \}, \quad k = 0, 1, \ldots, k_0.
\]

Subdivide the difference \( C_{k-1} \setminus C_k \) into smaller cubes with edges of length \( 2d = (1 - q)^{k-1} q \cdot d_0 \) such that the coordinates of their centers \( \xi \) are odd multiples of \( d \), and

\[
    |\xi| = (1 - q)^{k-1} (1 - \frac{q}{2}) \cdot d_0.
\]

For each value of \( k \) there are less than \( (2M)^n \) such smaller cubes, so there are \( N < k_0 \cdot (2M)^n \) cubes contained in \( C_0 \setminus C_{k_0} \); number them \( C_{(1)}, \ldots, C_{(N)} \).

Now let \( \epsilon > 0 \) be given. We shall show that for a proper choice of \( M \) and of \( k_0 = k_0(n) \), the right-hand side of

\[
    P \left( \sup_{\mu \in C_0} Z_n(\mu, 0) \geq 2\epsilon \right) \leq P \left( \sup_{\mu \in C_{k_0}} Z_n(\mu, 0) \geq 2\epsilon \right) + \sum_{i=1}^N P \left( \sup_{\mu \in C_{(i)}} Z_n(\mu, 0) \geq 2\epsilon \right)
\]

(51)

tends to 0 with increasing \( n \), which establishes the final result.

Actually, we shall choose

\[
    M \geq (3b)/(\epsilon a), \quad \text{which implies} \quad q \leq (\epsilon a)/(3b),
\]

(52)

and \( k_0 = k_0(n) \) is defined by

\[
    (1 - q)^{k_0} \cdot d_0 \leq n^{-\gamma'} < (1 - q)^{k_0-1} \cdot d_0
\]

(53)

where \( \frac{1}{2} < \gamma < 1 \) is an arbitrary fixed number. Thus, we have

\[
    k_0(n) - 1 < \frac{\gamma' \cdot (\log n + \log d_0)}{\log(1 - q)} \leq k_0(n),
\]

(54)

hence

\[
    N \leq \left( \frac{\gamma' \cdot (\log n + \log d_0)}{\log(1 - q)} + 1 \right) \cdot (2M)^n.
\]

(55)

Now take any of the cubes \( C_{(i)} \), with center \( \xi \) and edges of length \( 2d \) according to \( 2d = (1 - q)^{k-1} q \cdot d_0 \) and \( |\xi| = (1 - q)^{k-1} (1 - \frac{q}{2}) \cdot d_0 \). For \( \mu \in C_{(i)} \), we have then by (N-3),

\[
    |\lambda(\mu)| \geq a \cdot \|\mu\| \geq a \cdot (1 - q)^k \cdot d_0
\]

(56)
and

\[ |\lambda(\mu) - \lambda(\xi)| \leq \mathbb{E}_\tilde{x}[u(\tilde{x}, \xi, d)] \leq bd \leq b(1-q)^k q \cdot d_0. \tag{57} \]

We have

\[ Z_n(\mu, 0) \leq Z_n(\mu, \xi) + \frac{\sum_{j=1}^{n} [\psi(x_j, \xi) - \psi(x_j, 0) - \lambda(\xi)]}{\sqrt{n + n|\lambda(\mu)|}}, \tag{58} \]

hence

\[ \sup_{\mu \in \mathcal{C}_0} Z_n(\mu, 0) \leq U_n + V_n \tag{59} \]

with

\[ U_n = \frac{\sum_{j=1}^{n} [u(x_j, \xi, d) + \mathbb{E}_\tilde{x}[u(\tilde{x}, \xi, d)]]}{na(1-q)^k \cdot d_0}, \tag{60} \]

and

\[ V_n = \frac{\sum_{j=1}^{n} [\psi(x_j, \xi) - \psi(x_j, 0) - \lambda(\xi)]}{na(1-q)^k \cdot d_0}. \tag{61} \]

Thus,

\[ P(U_n \geq \epsilon) = P\left( \sum_{j=1}^{n} [u(x_j, \xi, d) - \mathbb{E}_\tilde{x}[u(\tilde{x}, \xi, d)]] \geq cna(1-q)^k \cdot d_0 - 2n\mathbb{E}_\tilde{x}[u(\tilde{x}, \xi, d)] \right). \tag{62} \]

In view of (57) and (52),

\[ ea(1-q)^k \cdot d_0 - 2\mathbb{E}_\tilde{x}[u(\tilde{x}, \xi, d)] \geq ea(1-q)^k \cdot d_0 - 2bq(1-q)^k \cdot d_0 \geq b(1-q)^k \cdot d_0, \tag{63} \]

hence (N-3) (iv) and Hoeffding’s inequality (Lemma 14) yields

\[ P(U_n \geq \epsilon) \leq \exp\left(-\frac{2bq^2(1-q)^{2k}d_0n}{c_2^2}\right) \leq \exp\left(-\frac{2bq^2(1-q)^{2k}d_0n}{c_2^2}\right) \leq \exp\left(-\frac{2bq^2(1-q)^{2k}d_0n}{c_2^2}\right). \tag{64} \]

In a similar way,

\[ P(V_n \geq \epsilon) \leq \frac{c}{gbq^2(1-q)^2} \cdot \frac{1}{n(1-q)^{k-1} \cdot d_0}. \tag{65} \]

Hence, we obtain from (53), (59), (64) and (65) that

\[ P\left( \sup_{\mu \in \mathcal{C}_0} Z_n(\mu, 0) \geq 2\epsilon \right) \leq n^{\gamma^*-1} \cdot \left( \frac{c}{b^2q(1-q)} + \frac{c}{gbq^2(1-q)^2} \right). \tag{66} \]

Furthermore,

\[ \sup_{\mu \in \mathcal{C}_0} Z_n(\mu, 0) \leq \frac{\sum_{j=1}^{n} [u(x_j, 0, d) + \mathbb{E}_\tilde{x}[u(\tilde{x}, 0, d)]]}{\sqrt{n}} \tag{67} \]

with \( d = (1-q)^k \cdot d_0 \leq n^{-\gamma^*} \). Hence,

\[ P\left( \sup_{\mu \in \mathcal{C}_0} Z_n(\mu, 0) \geq 2\epsilon \right) \leq P\left( \sum_{j=1}^{n} [u(x_j, 0, d) - \mathbb{E}_\tilde{x}[u(\tilde{x}, 0, d)]] \geq 2\sqrt{n}\epsilon - 2n\mathbb{E}_\tilde{x}[u(\tilde{x}, 0, d)] \right). \tag{68} \]

Since \( \mathbb{E}_\tilde{x}[u(\tilde{x}, 0, d)] \leq bd \leq bn^{-\gamma^*} \), set \( n_0 \) such that \( n_0^{\gamma^* - \frac{1}{2}} = \frac{2a}{\epsilon} \). Then, for \( n \geq n_0 \), we have

\[ \mathbb{E}_\tilde{x}[u(\tilde{x}, 0, d)] \leq \frac{\epsilon}{2\sqrt{n}} \Rightarrow 2\sqrt{n}\epsilon - 2n\mathbb{E}_\tilde{x}[u(\tilde{x}, 0, d)] \geq \sqrt{n}\epsilon; \]

thus, by Chebyshev’s inequality,

\[ P\left( \sup_{\mu \in \mathcal{C}_0} Z_n(\mu, 0) \geq 2\epsilon \right) \leq c \cdot \epsilon^{-2} \cdot n^{-\gamma^*}. \tag{69} \]

Now, putting (51), (55), (66), and (69) together, we obtain

\[ P\left( \sup_{\mu \in \mathcal{C}_0} Z_n(\mu, 0) \geq 2\epsilon \right) \leq c \cdot \epsilon^{-2} \cdot n^{-\gamma^*} + n^{\gamma^* - 1} \cdot \left( \frac{c}{b^2q(1-q)} + \frac{c}{gbq^2(1-q)^2} \right) \cdot \frac{(\gamma^* \cdot (\log n + \log d_0) + 1) \cdot (2M)^n}{|\log(1-q)|} \]

as long as \( n \geq n_0 \) where \( n_0 \) satisfying \( n_0^{\gamma^* - \frac{1}{2}} = \frac{2a}{\epsilon} \), which completes our proof of the lemma. \( \square \)
Appendix D: Missing Proofs for Section 3

Proof of Lemma 3. Our proof can be classified into three steps. We fix \( \{ \hat{q}_j^* \}_{j=1}^n \) as an optimal solution to (13). In the first step, we discretize both the convex optimization problem (12) and (13) into two LPs with a granularity \( K \) such that \( \{ \hat{q}_j^* \}_{j=1}^n \) is an optimal solution to the discretized LP of (13). In the second step, we show that we can select one optimal solution to the discretized LP of (12), such that the gap between the selected optimal solution and \( \{ \hat{q}_j^* \}_{j=1}^n \) can be bounded by a constant independent of the granularity \( K \). In the final step, we show that as the granularity \( K \) grows to infinity, there exists a subsequence of \( K \) such that the selected optimal solution to the discretized LP of (12) will converge to an optimal solution of (12), which completes our proof.

We now do the first step to discretize the convex optimization problem (12) and (13). For each \( j \in [n] \), we denote by a function

\[
G_j(q) = \int_{q'=1-q}^{1} F_{j}^{-1}(q')dq'.
\]

For any integer \( K \geq n + 1 \), we denote by a set \( \{ 0, \frac{1}{K-n}, \frac{2}{K-n}, \ldots, 1 \} \cup \{ \hat{q}_j^* \}_{j=1}^n \) and let \( q_k^K \) to be the \( k \)-th smallest element in this set, for \( k = 1, 2, \ldots, K + 1 \). For each \( j \in [n] \), we denote by \( \hat{G}_j^K(\cdot) \) the piece-wise linear interpolation of \( G_j \) based on the values at points \( \{ q_k^K \}_{k=1}^{K+1} \). From the concavity of the function \( G_j(\cdot) \), it is clear that we have

\[
\hat{G}_j^K(q) \leq G_j(q) \leq \hat{G}_j^K(q) + \frac{\max_{u \in [n]} \{ u_j \}}{K - n}, \quad \forall q \in [0, 1], \quad \text{and} \quad \hat{G}_j^K(q_j^*) = G_j(q_j^*). \tag{70}
\]

Then, we know that \( \{ \hat{q}_j^* \}_{j=1}^n \) is an optimal solution to the following optimization problem

\[
\hat{V}_{t,e}^K = \max \sum_{j=1}^{n} \left( \sum_{i=1}^{K} \beta_{j,k} \cdot x_{k,j} \right) \tag{71}
\]

subject to

\[
\sum_{j=1}^{n} \left( \sum_{i=1}^{K} \beta_{j,k} \cdot x_{k,j} \right) \leq c_i, \quad \forall i \in [m]
\]

\[
0 \leq x_{k,j} \leq q_{k+1} - q_k, \quad \forall j \in [n], \forall k \in [K]
\]

because optimization problem (71) differs from (13) only by having a pointwise-dominated objective function, and \( \{ \hat{q}_j^* \}_{j=1}^n \) attains the optimal objective value from (13) in the dominated problem (71). Without loss of generality, we assume that for each \( j \in [n] \), the piece-wise linear functions \( \hat{G}_j^K(\cdot) \) share the same set of end points of the piece-wise linear intervals, and we denote by \( \{ q_k^K \}_{k=1}^{K+1} \) the set of end points, with \( q_1^K = 0 \) and \( q_{K+1} = 1 \). Then, we have the following linear programming as a re-formulation of the discretization \( \hat{V}_{t,e}^K \) of the convex problem \( \hat{V}_{t,e}^{\text{Fid}} \) (13):

\[
\hat{V}_{t,e}^K = \max \sum_{j=1}^{n} \left( \sum_{i=1}^{K} \beta_{j,k} \cdot x_{k,j} \right) \tag{72}
\]

subject to

\[
\sum_{j=1}^{n} \left( \sum_{i=1}^{K} \beta_{j,k} \cdot x_{k,j} \right) \leq c_i, \quad \forall i \in [m]
\]

\[
0 \leq x_{k,j} \leq q_{k+1} - q_k, \quad \forall j \in [n], \forall k \in [K]
\]

where \( \beta_{j,k} \) is the coefficient that is inherited from the piece-wise linear function \( \hat{G}_j^K \). Here, the variable \( x_{k,j} \) can be interpreted as the number of queries, with type \( j \) and reward realization quantile lying in the interval
We do the third step to complete our proof. For any granularity $K$, our result is formalized in the following claim. Our analysis follows the analysis in Mangasarian and Shiau (1987) over the Lipschitz continuity of solutions of linear programming, and we further show an equivalence between two linear systems to obtain a bound that is independent of the granularity $K$.

**Claim 2.** There exists a constant $\mu$ such that for any optimal solution $\{\hat{x}_{k,j}^K\}$ of $\hat{V}_{t,e}^K(d)$, we can select one optimal solution $\{\bar{x}_{k,j}^K\}$ of $\bar{V}_e^K(d)$ such that

$$\left| \sum_{k=1}^{K} \bar{x}_{k,j}^K - \hat{x}_{k,j}^K \right| \leq \mu \cdot \max_{j'} \{ |p_{j'} \cdot s - d_{j'}| \}, \quad \forall j \in [n].$$

(74)

where

$$\mu = \sup_{v^1 \in \mathbb{R}^m, v^2, v^3 \in \mathbb{R}^n} \left\{ \left| \frac{v^1}{\|v^1\|} \right|, \left| \frac{v^2}{\|v^2\|} \right|, \left| \frac{v^3}{\|v^3\|} \right| \right\}$$

(75)

$$A = (a_{i,j})_{i \in [m], j \in [n]} \in \mathbb{R}^{m \times n} \text{ and } I_n \text{ is an identity matrix with a size } n \times n.$$

We do the third step to complete our proof. For any granularity $K$, let $\{\hat{x}_{k,j}^K\}$ be a solution of $\hat{V}_{t,e}^K(d)$ satisfying

$$\hat{x}_{k,j}^K = \begin{cases} p_{j} \cdot s \cdot (q_{k+1}^K - q_k^K), & \text{if } q_k^K \geq 1 - \hat{q}_j, \\ (q_{k+1}^K - (1 - \hat{q}_j^*) \cdot p_{j} \cdot s, & \text{if } q_k^K < 1 - \hat{q}_j \leq q_{k+1}^K \\ 0, & \text{if } q_{k+1}^K < 1 - \hat{q}_j. \end{cases}$$

Since $\{\hat{q}_j\}_{j=1}^n$ is an optimal solution to $\hat{V}_{t,e}^K$ under the formulation (71), we must have $\{\hat{x}_{k,j}^K\}$ is an optimal solution of $\hat{V}_{t,e}^K$, under the formulation (72). Then, we denote by $\{\bar{x}_{k,j}^K\}$ one optimal solution of $\bar{V}_e^K(d)$, as specified in Claim 2. We further construct

$$\hat{q}_j^K = \frac{\sum_{k=1}^{K} \bar{x}_{k,j}^K}{d_j}, \quad \forall j \in [n].$$

(76)

We denote by $\hat{q}^* = (\hat{q}_1^*, \ldots, \hat{q}_n^*)$ and $\hat{q}^K = (\hat{q}_1^K, \ldots, \hat{q}_n^K)$. From definition, we know that $\hat{q}^*, \hat{q}^K \in [0, 1]^K$. Therefore, there exists a point $\hat{q}^*$ and a sequence of integers $\{K_1, \ldots, K_w, \ldots\}_{w=1,2,\ldots}$ such that

$$\hat{q}^* = \lim_{w \to \infty} \hat{q}^{K_w}. \quad (77)$$

We show in the following claim that $\hat{q}^*$ is an optimal solution to $\bar{V}_e^{\text{Semi}}(d)$ (12).
Claim 3. Let $\bar{q}'$ be constructed in (77). Then, $\bar{q}'$ is an optimal solution to $\bar{V}_K^e(d)$ (12).

For any integer $w$, from Claim 2, we know that

$$
\|\bar{q}' - q^K_w\|_\infty = \left\| \sum_{j=1}^{K_w} \frac{x_j}{p_j \cdot s} - \sum_{j=1}^{K_w} \frac{x_j}{d_j} \right\|_\infty \\
\leq \max_{j \in [n]} \left\{ \frac{\mu}{p_j} \right\} \cdot \max_{j' \in [n]} \left\{ \left( \frac{d_j}{p_j \cdot s} - 1 \right) \right\}
$$

Therefore, from (77) and Claim 3, there exists an optimal solution $\bar{q}'$ of $\bar{V}_K^e(d)$ (12) such that

$$
\|\bar{q}' - q'\|_\infty \leq \max_{j \in [n]} \left\{ \frac{\mu}{p_j} \right\} \cdot \max_{j' \in [n]} \left\{ \left( \frac{d_j}{p_j \cdot s} - 1 \right) \right\}.
$$

Our proof is thus completed.

Proof of Claim 2. We denote the linear programming $\hat{V}_K^e(72)$ by

$$
\hat{V}_K^e = \max h^T x \\
s.t. \ \ l' \leq Ax \leq u'
$$

and denote the linear programming $\bar{V}_K^e(d)$ (73) by

$$
\bar{V}_K^e(d) = \max h^T x \\
s.t. \ \ l'' \leq Ax \leq u''
$$

Fix an optimal solution $\hat{x}^K$ of $\hat{V}_K^e$. We denote by $J_1, J_2$ and $J_3$ three row index sets of $A$ such that

$$
A_{J_1} \hat{x}^K = u'_{J_1}, A_{J_2} \hat{x}^K = l'_2, \text{ and } u'_{J_3} A_{J_3} \hat{x}^K < b'_3.
$$

We further fix an optimal solution $\bar{x}^K'$ of $\bar{V}_K^e(d)$ and denote by $J_1 = J_{1,1} \cup J_{1,2}$, $J_2 = J_{2,1} \cup J_{2,2}$ such that

$$
A_{J_1} \bar{x}^K' = u''_{J_1}, A_{J_2} \bar{x}^K' = l''_{J_2} \text{ and } A_{J_3} \bar{x}^K' = l''_3.
$$

Then, we denote by a set of linear equalities and linear inequalities

$$
A_{J_1} x = u''_{J_1},
A_{J_2} \bar{x}^K' \leq A_{J_2} x, A_{J_2} x \leq u''_{J_2},
A_{J_2} x = l''_{J_2},
A_{J_2} \bar{x}^K' \geq A_{J_2} x, A_{J_2} x \geq l''_{J_2},
l''_3 \leq A_{J_3} x \leq u''_{J_3}
$$

It is clear that $\bar{x}^K'$ satisfies the linear system (80). On the other hand, $\hat{x}^K$ satisfies the following linear system:

$$
A_{J_1} x = u'_{J_1},
A_{J_1} x - u''_{J_1} + A_{J_1} \bar{x}^K' \leq A_{J_1} x, A_{J_1} x \leq u'_{J_1},
A_{J_2} x = l'_{J_2},
l''_{J_2} - l''_{J_2} + A_{J_2} \bar{x}^K' \geq A_{J_2} x, A_{J_2} x \geq l''_{J_2},
l''_3 \leq A_{J_3} x \leq u''_{J_3}
$$
Our remaining analysis can be classified into two steps. For the first step, we show that any variable \( x \) that satisfies the linear system (80) turns out to be an optimal solution of \( \hat{V}_e^K(d) \) (79). For the second step, we show that for \( \hat{x}^K \) that satisfies the linear system (81), we can find a variable \( \hat{x}^K \) satisfying the linear system (80) such that (74) holds.

We now prove the first step. Since \( \hat{x}^K \) is an optimal solution of \( \hat{V}_e^K \) (78), from the KKT optimality condition, we know that there exists dual variables \( \mu_{j_1} \geq 0, v_{j_2} \leq 0 \) such that

\[
\begin{bmatrix} A_{j_1} \mu_{j_1} + A_{j_2} v_{j_2} = h \end{bmatrix}
\]

Then, for any \( x \) that satisfies the linear system (80), we have

\[
\begin{align*}
\hat{h}^T x &= \mu_{j_1} A_{j_1} x + v_{j_2} A_{j_2} x = \mu_{j_1} A_{j_1} x + \mu_{j_1,2} A_{j_1,2} x + v_{j_2} A_{j_2,1} x + v_{j_2,2} A_{j_2,2} x \\
&\geq \mu_{j_1,1} A_{j_1,1} \bar{x}^{K'} + \mu_{j_1,2} A_{j_1,2} \bar{x}^{K'} + v_{j_2,1} A_{j_2,1} \bar{x}^{K'} + v_{j_2,2} A_{j_2,2} \bar{x}^{K'} = \mu_{j_1} A_{j_1,1} \bar{x}^{K'} + v_{j_2} A_{j_2,2} \bar{x}^{K'} \\
&= \hat{h}^T \bar{x}^{K'} = \hat{V}_e^K(d)
\end{align*}
\]

where the first inequality follows from both \( x, \bar{x}^{K'} \) satisfies the linear system (80) and thus \( A_{j_1,1} x = A_{j_1,1} \bar{x}^{K'} = u_{j_1,1}', A_{j_2,1} x = A_{j_2,1} \bar{x}^{K'} = l_{j_2,1}', A_{j_1,2} \bar{x}^{K'} \leq A_{j_1,2} x \) and \( A_{j_2,2} \bar{x}^{K'} \geq A_{j_2,2} x \).

It only remains to show the second step. We prove by exploiting the special structure of the linear systems (80) and (81). Note that under our setting, we have

\[
A = \begin{bmatrix} A & \ldots & A \\ I_{nK} \end{bmatrix}
\]

where \( I_{nK} \) denotes an identity matrix with size \( nK \times nK \) and \( A = (a_{i,j})_{i \in [m], j \in [n]} \in \mathbb{R}^{m \times n} \). Also, the linear system (80) can be rewritten as

\[
l^1 \leq Ax \leq u^1
\]

and the linear system (81) can be rewritten as

\[
l^2 \leq Ax \leq u^2.
\]

Then, for the solution \( \hat{x}^K \) of the linear system (83), we construct \( \hat{y} \in \mathbb{R}^n \) with

\[
\hat{y}_j = \sum_{k=1}^{K} \hat{x}_{k,j}^K, \quad \forall j \in [n].
\]

It is clear that \( \hat{y} \) is a solution to the following linear system

\[
\hat{l}_2 \leq \begin{bmatrix} A \\ I_n \end{bmatrix} y \leq \hat{u}_2
\]

where \( \hat{l}_2 \in \mathbb{R}^{n+m} \) satisfying

\[
\hat{l}_2^i = l_i^2, \quad \forall i \in [m] \quad \text{and} \quad \hat{l}_2^j \equiv \frac{K}{k=1} l_{j+k,1}, \quad \forall j \in [n]
\]

and \( \hat{u}_2 \in \mathbb{R}^{n+m} \) satisfying

\[
\hat{u}_2^i = u_i^2, \quad \forall i \in [m] \quad \text{and} \quad \hat{u}_2^j \equiv \frac{K}{k=1} u_{j+k,1}, \quad \forall j \in [n].
\]
In the same way, we denote by \( \hat{l}_i^1 \in \mathbb{R}^{n+m} \) satisfying
\[
\hat{l}_i^1 = l_i^1, \quad \forall i \in [m] \quad \text{and} \quad \hat{l}_{n+j}^1 = \sum_{k=1}^{K} l_{j+kn}^1, \quad \forall j \in [n]
\]
and \( \hat{u}_i^1 \in \mathbb{R}^{n+m} \) satisfying
\[
\hat{u}_i^1 = u_i^1, \quad \forall i \in [m] \quad \text{and} \quad \hat{u}_{n+j}^1 = \sum_{k=1}^{K} u_{j+kn}^1, \quad \forall j \in [n].
\]
We consider the linear system
\[
\hat{l}^1 \leq \begin{bmatrix} A \\ I_n \end{bmatrix} y \leq \hat{u}^1 \tag{85}
\]
From Lemma 11, we know that for the solution \( \hat{y} \) that satisfies linear system (84), there exists a solution \( \hat{y} \) that satisfies linear system (85) and there also exists a constant \( \mu \) such that
\[
\| \hat{y} - \bar{y} \|_{\infty} \leq \mu \cdot \| (\hat{l}^1, \hat{u}^1) - (\hat{l}, \hat{u}) \|_{\infty} \leq \mu \cdot \max_{j' \in [n]} \{ (p_{j'} \cdot s - d_{j'}) \}. \tag{86}
\]
Note that here the constant \( \mu \) depends solely on \( A \) and \( I_n \) and is independent of the granularity \( K \). We now construct a solution to the linear system (82) from \( \hat{y} \) to complete our proof. For each \( j \in [n] \) and each \( k \in [K] \), we define
\[
\bar{x}_{k,j}^K = l_{j+kn}^1 + (\bar{y}_j - \hat{l}_{n+j}^1) \cdot \frac{u_{n+j}^1 - l_{n+j}^1}{\hat{u}_{n+j}^1 - \hat{l}_{n+j}^1}.
\]
It is clear to see that \( \bar{x}^K \) satisfies linear system (82) and satisfies
\[
\sum_{k=1}^{K} \bar{x}_{k,j}^K = \bar{y}_j, \quad \forall j \in [n]. \tag{87}
\]
Our proof of the second step is completed from (86), (77) and the fact that \( \bar{x}^K \) satisfies linear system (82).
From the conclusion of the first step, we know that \( \bar{x}^K \) is an optimal solution to \( V_e^K(d) \) and our proof of Claim 2 is thus completed.

**Proof of Claim 3.** For any integer \( w \), it is clear to see that \( \bar{q}^K_w \) is a feasible solution to \( V_e^{Semi}(d) \) (12), which implies that the limiting point \( \bar{q}^* \) is a feasible solution to \( V_e^{Semi}(d) \) (12) since the feasible set is closed.
We now prove optimality.

It is direct to see from the construction of the piece-wise linear function \( \hat{G}_j^K(\cdot) \) that
\[
\hat{V}_e^{K_w}(d) = \sum_{j=1}^{n} \sum_{k=1}^{K_w} \beta_{j,k} \cdot \bar{x}_{k,j} \leq \sum_{j=1}^{n} d_j \cdot \int_{q_{j}^w}^{1} F_j^{-1}(q) dq \leq \hat{V}_e^{Semi}(d).
\]
where the first inequality follows from the definition \( \hat{q}^{K_w}_j = \frac{\sum_{k=1}^{K_w} \bar{x}_{k,j}}{p_{j} \cdot s} \).

We now show that \( \hat{V}_e^{Semi}(d) \leq \hat{V}_e^{K_w}(d) + \frac{(\sum_{j=1}^{n} d_j) \cdot \max_{j \in [n]}(u_{j})}{K_w} \). For an optimal solution \( \{ \bar{q}_{j}^* \}_{j=1}^{n} \) of \( V_e^{Semi}(d) \).
We construct for any \( k \in [K_w] \) and \( j \in [n] \)
\[
\hat{x}_{k,j}^{K_w} = \begin{cases} p_{j} \cdot s \cdot (\bar{q}^{K_w}_{k+1} - \bar{q}^{K_w}_{k}), & \text{if } \bar{q}^{K_w}_{k} \geq 1 - \bar{q}_{j}^* \\ (\bar{q}^{K_w}_{k+1} - (1 - \bar{q}_{j}^*)) \cdot p_{j} \cdot s, & \text{if } \bar{q}^{K_w}_{k} < 1 - \bar{q}_{j}^* \leq \bar{q}^{K_w}_{k+1} \\ 0, & \text{if } \bar{q}^{K_w}_{k+1} < 1 - \bar{q}_{j}^*. \end{cases}
\]
It is clear to see that \( \{ x^K_{k,j} \} \) is a feasible solution to \( \bar{V}^K_e(d) \) (73) and it holds that

\[
\bar{V}^{\text{Semi}}_e(d) = \sum_{j=1}^n d_j \cdot G_j(q_j') \leq \sum_{j=1}^n d_j \cdot \hat{G}_j^K(q_j') + \left( \sum_{j=1}^n d_j \cdot \max_{j \in [n]} \{ u_j \} \right) / K_w - n = \sum_{j=1}^n \sum_{k=1}^{K_w} d_j \cdot x^K_{k,j} + \left( \sum_{j=1}^n d_j \cdot \max_{j \in [n]} \{ u_j \} \right) / K_w - n \leq \bar{V}^K_e(d) + \left( \sum_{j=1}^n d_j \cdot \max_{j \in [n]} \{ u_j \} \right) / K_w - n
\]

where the first inequality follows from (70). Therefore, we conclude that

\[
\bar{V}^K_e(d) \leq \sum_{j=1}^n p_j \cdot s \cdot \int_{q=1-q^K_w}^1 F_j^{-1}(q) dq \leq \bar{V}^{\text{Semi}}_e(d) \leq \bar{V}^K_e(d) + \left( \sum_{j=1}^n d_j \cdot \max_{j \in [n]} \{ u_j \} \right) / K_w - n.
\]

Note that

\[
\sum_{j=1}^n d_j \cdot \int_{q=1-q^K_w}^1 F_j^{-1}(q) dq = \lim_{w \to \infty} \sum_{j=1}^n d_j \cdot \int_{q=1-q^K_w}^1 F_j^{-1}(q) dq.
\]

We have

\[
\left| \sum_{j=1}^n d_j \cdot \int_{q=1-q^K_w}^1 F_j^{-1}(q) dq - \bar{V}^{\text{Semi}}_e(d) \right| = \lim_{w \to \infty} \left| \sum_{j=1}^n d_j \cdot \int_{q=1-q^K_w}^1 F_j^{-1}(q) dq - \bar{V}^{\text{Semi}}_e(d) \right| \leq \lim_{w \to \infty} \left( \sum_{j=1}^n d_j \cdot \max_{j \in [n]} \{ u_j \} \right) / K_w - n = 0
\]

which implies that \( q' \) is an optimal solution to \( \bar{V}^{\text{Semi}}_e(d) \) (12). Our proof is therefore completed. \( \square \)

**Proof of Claim 1.** Denote by functions

\[
G_j(q) := \int_{q'=q}^1 F_j^{-1}(q') dq', \quad \forall j \in [n].
\]

It is clear that

\[
G'_j(q) = -F_j^{-1}(q) \quad \text{and} \quad G''_j(q) = -\frac{1}{f(q|a_j)} \in [-\frac{1}{\alpha}, 0].
\]

Therefore, from the concavity of \( G_j(\cdot) \), we have

\[
\int_{q=q_1}^{q_2} F_j^{-1}(q) dq = G_j(q_1) - G_j(q_2) \leq F_j^{-1}(q_1) \cdot (q_2 - q_1) + \frac{(q_2 - q_1)^2}{\alpha}
\]

which completes our proof. \( \square \)

**Proof of Lemma 4.** The formulation of \( \mu \) in (75) is copied as follows

\[
\mu = \sup_{v^1 \in \mathbb{R}^n, v^2, v^3 \in \mathbb{R}^n} \left\{ \begin{array}{l} \|v^1\|_1 + \|v^2\|_1 + \|v^3\|_1 \leq \|v^1\|_1 + \|v^2\|_1 + \|v^3\|_1 = 1, \\
\text{Rows of } \begin{bmatrix} A \\ I_n \end{bmatrix} \text{ corresponding to non-zero elements} \\
\text{of } \begin{pmatrix} v^1 \\ v^2 \\ v^3 \end{pmatrix} \text{ are linear independent.} \end{array} \right\}
\]

Here, \( A = (a_{i,j})_{i \in [n], j \in [n]} \in \mathbb{R}^{n \times n} \) and \( I_n \) is an identity matrix with a size \( n \times n \). From Theorem 2.2 of Mangasarian and Shiau (1987), we know that \( \mu \) in (88) is finite. We now derive an upper bound of it.

From Proposition 2.6 of Mangasarian and Shiau (1987), we know that

\[
\mu \leq (m + 2n) \cdot \Delta A
\]
where $\Delta A$ denotes the maximum of the absolute values of the determinants of the square submatrices of
\[
\begin{bmatrix}
A & I_n \\
I_n & -I_n
\end{bmatrix}.
\]

We denote by $A'$ to be a square submatrix of
\[
\begin{bmatrix}
A & I_n \\
I_n & -I_n
\end{bmatrix}.
\] Following the Hadamard’s inequality on matrix determinant (Lemma 12), we have that
\[
|\det(A')| = \prod_{i\in|n'|} \|a'_i\|_2
\]
where $n'$ denotes the size of matrix $A'$ and $a'_i$ denotes the $i$-th column of matrix $A'$. Therefore, it is clear to see that
\[
|\det(A')| = \prod_{i\in|n'|} \|a'_i\|_2 \leq \bar{a}^{n'} \leq \bar{a}^n
\]
for any submatrix $A'$, which completes our proof. □

Appendix E: Missing Proofs for Section 4

Proof of Lemma 5. By plugging (37) into (8), we have that
\[
\text{Myopic}_i(\pi, \hat{e}_t^i) = E_{\hat{r}_t, \hat{a}_t} \left[ E_{I_{t+1}} \left[ V_{\hat{e}_t}^\text{Off} (I_t) \right] - V_{\hat{e}_t}^\text{Off} (\tilde{I}_t) (I_{t+1}) \right] - \hat{r}_t \cdot \hat{x}_t^i
\]
\[
= E_{\hat{r}_t, \hat{a}_t} \left[ E_{I_{t+1}} \left[ \hat{x}_t^{\text{round}} - \hat{x}_t^i \right] \cdot \left( \hat{r}_t + V_{\hat{e}_t}^\text{Off} (I_{t+1}) + (\hat{x}_t^i - \hat{x}_t^{\text{round}}) \cdot \hat{r}_t - M_{\hat{e}_t} \hat{a}_t (I_{t+1}) \right) \right]
\]
\[
= \int_{\hat{r}_t, \hat{a}_t} \left[ E_{I_{t+1}} \left[ \hat{x}_t^{\text{round}} - \hat{x}_t^i \right] \cdot \left( \hat{r}_t - M_{\hat{e}_t} \hat{a}_t (I_{t+1}) \right) \right] E[\hat{r}_t, \hat{a}_t]
\]
\[
\leq \tilde{a} \cdot E_{I_{t+1}} \left[ E_{\hat{r}_t, \hat{a}_t} \left[ \left( \hat{r}_t - M_{\hat{e}_t} \hat{a}_t (I_{t+1}) \right)^2 \right] \cdot \mathbb{I} (\hat{e}_t \geq \tilde{a}) \right]
\]
\[
\leq 2 \tilde{a} \cdot \int_{\hat{e}_t \geq \tilde{a}} E_{\hat{a}_t} \left[ \left( \hat{e}_t - M_{\hat{e}_t} \hat{a}_t (I_{t+1}) \right)^2 \right]
\]
where the second inequality follows from Assumption 3 which implies that the density of $\hat{r}_t$ is upper bounded by $\tilde{a}$ and the third inequality follows from $(a+b)^2 \leq 2a^2 + 2b^2$ for any $a, b$. Plugging (91) into (90), we complete our proof of the lemma. □

Proof of Lemma 7. Note that we have
\[
L_{c, t+1}^\text{FL} (\mu) = \frac{c^\top \mu}{s - 1} + \int_{\tilde{a} \geq \tilde{a}} \left[ E_{\hat{r} \sim F_{\hat{a}}} (\hat{r} - \tilde{a}^\top \mu) \right] \mathbb{I} (\hat{r} \geq \tilde{a}) f(r) dr
\]
which implies that
\[
\frac{\partial L_{c, t+1}^\text{FL} (\mu)}{\partial \mu_i} = \frac{c_i}{s - 1} - E_{\hat{a}} [\hat{a}_i \cdot (1 - F(\hat{a}^\top \mu | \hat{a}))], \quad \forall i \in [m]
\]
and
\[
\frac{\partial^2 L_{c,t+1}^{\text{Fld}}(\mu)}{\partial \mu_i \partial \mu_i'} = E_a[f(\tilde{a}^\top \mu | \tilde{a}) \cdot \tilde{a}_i \cdot \tilde{a}_i'], \quad \forall i, i' \in [m]
\]
where we use \( \tilde{a} \) to denote the \( i \)-th element of vector \( \tilde{a} \). Moreover, note that for any \( \mu \geq 0 \), from Taylor’s theorem, there exists a \( \mu' \) that lies on the intersection between \( \mu \) and \( \mu^* \) such that
\[
L_{c,t+1}^{\text{Fld}}(\mu) = L_{c,t+1}^{\text{Fld}}(\mu^*) + \frac{1}{2} \cdot (\mu - \mu^*)^\top \frac{\partial^2 L_{c,t+1}^{\text{Fld}}(\mu^*)}{\partial \mu^2} \cdot (\mu - \mu^*)
\]
where \( \frac{\partial L_{c,t+1}^{\text{Fld}}(\mu^*)}{\partial \mu_i} = \left( \frac{\partial L_{c,t+1}^{\text{Fld}}(\mu^*)}{\partial \mu_1}, \ldots, \frac{\partial L_{c,t+1}^{\text{Fld}}(\mu^*)}{\partial \mu_m} \right) \), and \( \frac{\partial^2 L_{c,t+1}^{\text{Fld}}(\mu^*)}{\partial \mu^2} \) is the Hessian matrix that equals
\[
E_a[f(\tilde{a}^\top \mu | \tilde{a}) \cdot \tilde{a} \cdot \tilde{a}^\top].
\]
Note that \( \Omega \) is assumed to be a convex set in Assumption 3 and both \( \mu, \mu^* \in \Omega \).

Then, we have \( \mu' \in \Omega \) and Assumption 3 implies that \( E_a[f(\tilde{a}^\top \mu | \tilde{a})] \geq \alpha \). The smallest positive eigenvalue of \( \frac{\partial^2 L_{c,t+1}^{\text{Fld}}(\mu^*)}{\partial \mu^2} \) is therefore lower bounded by \( \alpha \cdot \beta \). Also, from the optimality of \( \mu^* \), we must have
\[
(\mu - \mu^*)^\top \frac{\partial L_{c,t+1}^{\text{Fld}}(\mu^*)}{\partial \mu} \geq 0
\]
for any \( \mu \geq 0 \). Thus, it holds that
\[
\frac{\alpha \beta}{2} \|P_3(\mu - \mu^*)\|_2^2 \leq \frac{1}{2} (\mu - \mu^*)^\top \frac{\partial^2 L_{c,t+1}^{\text{Fld}}(\mu^*)}{\partial \mu^2} \cdot (\mu - \mu^*) \leq L_{c,t+1}^{\text{Fld}}(\mu) - L_{c,t+1}^{\text{Fld}}(\mu^*)
\]
which completes our proof. \( \square \)

**Proof of Lemma 8.** From (44), we know that
\[
M_{c,a_i}(I_{t+1}) - E[M_{c,a_i}(I_{t+1})] \leq a_i^\top \mu_2 - E[a_i^\top \tilde{\mu}_1], \quad \text{if } M_{c,a_i}(I_{t+1}) \geq E[M_{c,a_i}(I_{t+1})]
\]
\[
E[M_{c,a_i}(I_{t+1})] - M_{c,a_i}(I_{t+1}) \leq E[a_i^\top \tilde{\mu}_2] - a_i^\top \mu_1, \quad \text{if } M_{c,a_i}(I_{t+1}) \leq E[M_{c,a_i}(I_{t+1})]
\]
which implies that
\[
\text{Var}(M_{c,a_i}(I_{t+1})) \leq 2\text{Var}(a_i^\top \tilde{\mu}_1) + 2\text{Var}(a_i^\top \tilde{\mu}_2) + 2(E[a_i^\top \tilde{\mu}_1] - E[a_i^\top \tilde{\mu}_2])^2
\]
Note that we have
\[
\text{Var}(a_i^\top \tilde{\mu}_1) = E[(a_i^\top \tilde{\mu}_1 - a_i^\top \mu_1 + a_i^\top \mu_1 - a_i^\top E[\tilde{\mu}_1])^2] \leq 2E[(a_i^\top \mu_1 - a_i^\top \mu_1)^2] + 2(a_i^\top \mu_1 - a_i^\top E[\tilde{\mu}_1])^2
\]
Similarly, we have
\[
\text{Var}(a_i^\top \tilde{\mu}_2) = E[(a_i^\top \tilde{\mu}_2 - a_i^\top \mu_2 + a_i^\top \mu_2 - a_i^\top E[\tilde{\mu}_2])^2] \leq 2E[(a_i^\top \mu_2 - a_i^\top \mu_2)^2] + 2(a_i^\top \mu_2 - a_i^\top E[\tilde{\mu}_2])^2
\]
Also, we have
\[
(E[a_i, \tilde{\mu}_2] - E[a_i, \tilde{\mu}_1])^2 = (a_i^\top E[\tilde{\mu}_2] - a_i^\top \mu_2 + a_i^\top \mu_2 - a_i^\top \nu_1 + a_i^\top \mu_1 - a_i^\top E[\tilde{\mu}_1])^2
\]
\[
\leq 3(a_i^\top E[\mu_2] - a_i^\top \mu_2)^2 + 3(a_i^\top \mu_2 - a_i^\top \mu_1)^2 + 3(a_i^\top \mu_1 - a_i^\top E[\tilde{\mu}_1])^2
\]
Thus, we have that
\[
\text{Var}(M_{c,a_i}(I_{t+1})) \leq 10(a_i^\top E[\tilde{\mu}_1] - a_i^\top \mu_1)^2 + 10(a_i^\top E[\tilde{\mu}_2] - a_i^\top \mu_2)^2 + 6(a_i^\top \mu_2 - a_i^\top E[\tilde{\mu}_2])^2
\]
\[
+ 4E[(a_i^\top \mu_1 - a_i^\top \mu_1)^2] + 4E[(a_i^\top \mu_2 - a_i^\top \mu_2)^2]
\]
\[
\leq 14E[(a_i^\top \mu_1 - a_i^\top \mu_1)^2] + 14E[(a_i^\top \mu_2 - a_i^\top \mu_2)^2] + 6(a_i^\top \mu_2 - a_i^\top \mu_1)^2.
\]

It only remains to bound the term \( (a_i^\top \mu_2 - a_i^\top \mu_1)^2 \). We have the following claim, where the proof is relegated to the end of this proof.

**Claim 4.** It holds that
\[
(a_i^\top (\mu_1 - \mu_2))^2 \leq \frac{2}{(8 - 1)\alpha \beta} \cdot \bar{d}^2 \cdot m^{1/2} \cdot \gamma
\]
where \( \bar{d} = \max_{j \in [n]} \|a_j\|_2 \) and \( \gamma = \max_{i \in [m], j \in [n]: a_j > 0} \frac{a_j^2}{a_{ij}} \).

Therefore, our proof is completed by combining (92) and Claim 4. \( \square \)
Proof of Claim 4. From Lemma 7, by substituting $\tilde{\mu}_1$ into $\mu^*$, and substituting $\tilde{\mu}_2$ into $\mu$, we have that
\[
\frac{\alpha \beta}{2} \cdot \|P(\tilde{\mu}_2 - \mu_1)\|^2 \leq L_{c,t+1}(\tilde{\mu}_2) - L_{c,t+1}(\mu_1) = L_{c,t+1}(\tilde{\mu}_2) - L_{c,t+1}(\tilde{\mu}_2) + L_{c-a_t,t+1}(\tilde{\mu}_2) - L_{c-a_t,t+1}(\mu_1)
\]
where the second inequality follows from $L_{c-a_t,t+1}(\tilde{\mu}_2) \leq L_{c-a_t,t+1}(\mu_1) \leq L_{c-a_t,t+1}(\tilde{\mu}_2)$ by noting that $\tilde{\mu}_2 \in \arg\min_{\mu \geq 0} L_{c-a_t,t+1}(\mu)$. Thus, we have
\[
(a_t^\top (\tilde{\mu}_2 - \mu_1))^2 \leq d^2 \cdot \|P(\tilde{\mu}_2 - \mu_1)\|^2 \leq \frac{2}{s-1} \cdot (a_t^\top \tilde{\mu}_2) \cdot d^2 \leq \frac{2}{s-1} \cdot \|a_t\|_2 \cdot \|\tilde{\mu}_2\|_2 \cdot d^2
\]
Clearly, we must have $\hat{\mu}_2 \leq \gamma$ for each $i \in [m]$ with $\gamma = \max_{i \in [m], j \in [n/a_j], o > 0} \frac{\mu_j}{\alpha_j}$, which implies that $\|\tilde{\mu}_2\|_2 \leq \sqrt{m} \gamma$. Thus, it holds that
\[
(a_t^\top (\tilde{\mu}_2 - \mu_1))^2 \leq \frac{2}{(s-1)\alpha \beta} \cdot d^3 \cdot m^{1/2} \cdot \gamma
\]
which completes our proof. \(\square\)

Proof of Lemma 9. We first consider the setting where $c \geq \tilde{a}_t$ and thus both $\tilde{\mu}_1$ and $\tilde{\mu}_2$ are well-defined in (43). We show that $\tilde{x}_t^* \neq \tilde{x}_t^{\text{round}}$ happens only if $a_t^\top \tilde{\mu}_1 \leq \tilde{r}_t \leq a_t^\top \tilde{\mu}_2$. Note that the value of $\tilde{x}_t^*$ can be determined in the following way:
\[
\tilde{x}_t^* = \arg\max_{\phi \in [0,1]} \phi \cdot \tilde{r}_t + V_{c-\phi \tilde{a}_t}(I_{t+1}) - V_c(I_{t+1})
\]
Note that for any $\phi \in [0,1]$, we must have
\[
-\phi \cdot a_t^\top \tilde{\mu}_2 \leq V_{c-\phi \tilde{a}_t}(I_{t+1}) - V_c(I_{t+1}) \leq -\phi \cdot a_t^\top \tilde{\mu}_1
\]
which can be proved following the same intuition of Lemma 6. Therefore, when $\tilde{r}_t \leq a_t^\top \tilde{\mu}_1$, we must have $\tilde{x}_t^* = \tilde{x}_t^{\text{round}} = 0$, and when $\tilde{r}_t \geq a_t^\top \tilde{\mu}_2$, we must have $\tilde{x}_t^* = \tilde{x}_t^{\text{round}} = 1$. We conclude that $G_c(I_t) = 0$ when $\tilde{r}_t \leq a_t^\top \tilde{\mu}_1$ or $\tilde{r}_t \geq a_t^\top \tilde{\mu}_2$. We now assume that $a_t^\top \tilde{\mu}_1 \leq \tilde{r}_t \leq a_t^\top \tilde{\mu}_2$ and we further consider two cases as follows:

Case 1: If $\tilde{x}_t^{\text{round}} = 0$, then we have
\[
G_c(I_t) = \tilde{r}_t \cdot \tilde{x}_t^* + V_{c-a_t}(I_{t+1}) - V_c(I_{t+1})
\]
Note that we have
\[
V_{c-a_t}(I_{t+1}) - V_c(I_{t+1}) \leq -\tilde{x}_t^* \cdot a_t^\top \tilde{\mu}_1
\]
By noting $\tilde{r}_t \leq a_t^\top \tilde{\mu}_2$ and $\tilde{x}_t^* \leq 1$, we have that
\[
G_c(I_t) \leq \tilde{x}_t^* \cdot (a_t^\top \tilde{\mu}_2 - a_t^\top \tilde{\mu}_1) \leq a_t^\top \tilde{\mu}_2 - a_t^\top \tilde{\mu}_1
\]
which gives us an upper bound for the first case.

Case 2: If $\tilde{x}_t^{\text{round}} = 1$, then we have
\[
G_c(I_t) = \tilde{r}_t \cdot (\tilde{x}_t^* - 1) + V_{c-a_t}(I_{t+1}) - V_c(I_{t+1})
\]
Note that we have
\[
V_{c-a_t}(I_{t+1}) - V_c(I_{t+1}) \leq (1 - \tilde{x}_t^*) \cdot a_t^\top \tilde{\mu}_2.
\]
By noting \( \tilde{r}_t \geq \tilde{a}_t^\top \tilde{\mu}_1 \) and \( 0 \leq 1 - \tilde{x}_t^* \leq 1 \), we have that

\[
G_c(I_t) \leq (1 - \tilde{x}_t^*) \cdot (\tilde{a}_t^\top \tilde{\mu}_2 - \tilde{a}_t^\top \tilde{\mu}_1) \leq \tilde{a}_t^\top \tilde{\mu}_2 - \tilde{a}_t^\top \tilde{\mu}_1
\]

which gives us an upper bound for the second case. Therefore, on both cases, we show that if \( \tilde{a}_t^\top \tilde{\mu}_1 \leq \tilde{r}_t \leq \tilde{a}_t^\top \tilde{\mu}_2 \), it holds that

\[
G_c(I_t) \leq \tilde{a}_t^\top \tilde{\mu}_2 - \tilde{a}_t^\top \tilde{\mu}_1.
\] (94)

which implies that

\[
\mathbb{E}_{\tilde{a}_t} \left[ \mathbb{E}_{t+1} \left[ \mathbb{E}_{r \sim F(\{\tilde{a}_t\})}[G_c(I_t)] \right] \right] \leq \mathbb{E}_{\tilde{a}_t} \left[ \mathbb{E}_{t+1} \left[ \mathbb{E}_{r \sim F(\{\tilde{a}_t\})}[\mathbb{I}_{\tilde{a}_t^\top \tilde{\mu}_1 \leq \tilde{r}_t \leq \tilde{a}_t^\top \tilde{\mu}_2 \cdot (\tilde{a}_t^\top \tilde{\mu}_2 - \tilde{a}_t^\top \tilde{\mu}_1)] \right] \right]
\]

\[
\leq \tilde{\alpha} \cdot \mathbb{E}_{\tilde{a}_t} \left[ \mathbb{E}_{t+1} \left[ (\tilde{a}_t^\top \tilde{\mu}_2 - \tilde{a}_t^\top \tilde{\mu}_1)^2 \right] \right].
\]

We further note that

\[
(\tilde{a}_t^\top \tilde{\mu}_2 - \tilde{a}_t^\top \tilde{\mu}_1)^2 = (\tilde{a}_t^\top \tilde{\mu}_2 - \tilde{a}_t^\top \tilde{\mu}_2 + \tilde{a}_t^\top \tilde{\mu}_2 - \tilde{a}_t^\top \tilde{\mu}_1 + \tilde{a}_t^\top \tilde{\mu}_1 - \tilde{a}_t^\top \tilde{\mu}_1)^2
\]

\[
\leq 3(\tilde{a}_t^\top \tilde{\mu}_2 - \tilde{a}_t^\top \tilde{\mu}_2)^2 + 3(\tilde{a}_t^\top \tilde{\mu}_2 - \tilde{a}_t^\top \tilde{\mu}_1)^2 + 3(\tilde{a}_t^\top \tilde{\mu}_1 - \tilde{a}_t^\top \tilde{\mu}_1)^2.
\]

We use Claim 4 to bound the term \( (\tilde{a}_t^\top \tilde{\mu}_2 - \tilde{a}_t^\top \tilde{\mu}_1)^2 \). Then, we have

\[
\mathbb{E}_{t+1} \left[ G_c(I_t) \right] \leq \frac{\kappa_2}{s-1} + \kappa_2 \cdot \mathbb{E}_{\tilde{a}_t} \left[ \mathbb{E}_{t+1} \left[ (\tilde{a}_t^\top \tilde{\mu}_1 - \tilde{\alpha}_{\tilde{a}_t} \tilde{\mu}_1)^2 \right] \right] + \kappa_2 \cdot \mathbb{E}_{\tilde{a}_t} \left[ \mathbb{E}_{t+1} \left[ (\tilde{a}_t^\top \tilde{\mu}_2 - \tilde{\alpha}_{\tilde{a}_t} \tilde{\mu}_2)^2 \right] \right]
\] (95)

for a constant \( \kappa_2 \).

We then consider the setting where \( c \geq \tilde{a}_t \) does not hold and only \( \tilde{\mu}_1 \) is well-defined in (43). Still, when \( \tilde{r}_t \leq \tilde{a}_t^\top \tilde{\mu}_1 \), we must have \( \tilde{x}_t^* = \tilde{x}_t^\text{umnd} = 0 \), which implies \( G_c(I) = 0 \). We now assume that \( \tilde{r}_t \geq \tilde{a}_t^\top \tilde{\mu}_1 \). Then, since

\[
-\phi \cdot \tilde{a}_t^\top \tilde{\mu}_2 \leq \tilde{v}_t^\text{Off} - \tilde{v}_t^\text{Off}(I_t+1) \leq -\phi \cdot \tilde{a}_t^\top \tilde{r}_t
\]

for any \( \phi \in [0,1] \), we have

\[
G_c(I_t) \leq \tilde{x}_t^* \cdot (\tilde{r}_t - \tilde{a}_t^\top \tilde{\mu}_1) \leq \tilde{r}_t - \tilde{a}_t^\top \tilde{\mu}_1
\]

by noting \( \tilde{x}_t^* \in [0,1] \). Denote by \( j_t \) as the type of query \( t \). It holds that

\[
\mathbb{E}_{r \sim F(\{\tilde{a}_t\})}[G_c(I_t)] \leq \frac{1}{2} \cdot (u_{j_t} - a_{j_t}^\top \tilde{\mu}_1)^2 \leq (u_{j_t} - a_{j_t}^\top \tilde{\mu}_1)^2 + (a_{j_t}^\top \tilde{\mu}_1 - a_{j_t}^\top \tilde{\mu}_1)^2 \leq u_{j_t} \cdot (u_{j_t} - a_{j_t}^\top \tilde{\mu}_1) + (a_{j_t}^\top \tilde{\mu}_1 - a_{j_t}^\top \tilde{\mu}_1)^2.
\]

It is clear to see that \( \min_{\alpha} \tilde{L}_{c+1}^\text{Fid}(\mu) \) is the dual problem of \( \max_{\bar{x}} \tilde{V}_{t+1}^\text{Fid}(\bar{x}) \). Denote by \( \bar{x} = \{ \bar{x}_j(r) \} \in \arg \max_{\bar{x}} \tilde{V}_{t+1}^\text{Fid}(\bar{x}) \) such that \( \bar{x} \) and \( \tilde{\mu}_1 \) is an optimal primal-dual pair. Then, when \( c < a_{j_t,i} \) for some \( i \in [m] \), we must have

\[
\mathbb{E}_{r \sim F(\{\tilde{a}_t\})}[\bar{x}_{j_t}(r)] \leq \frac{1}{p_{j_t} \cdot (s-1)}
\]

and as a result we have

\[
\alpha \cdot (u_{j_t} - a_{j_t}^\top \tilde{\mu}_1) \leq P(a_{j_t}^\top \tilde{\mu}_1 \leq r \leq u_{j_t} | r \sim F_{j_t}) = \mathbb{E}_{r \sim F(\{\tilde{a}_t\})}[\bar{x}_{j_t}(r)] \leq \frac{1}{p_{j_t} \cdot (s-1)}
\]

where \( \alpha > 0 \) is a lower bound on the density function specified in Assumption 1. Then we have

\[
\mathbb{E}_{t+1} \left[ \mathbb{E}_{r \sim F(\{\tilde{a}_t\})}[G_c(I_t)] \right] \leq \frac{u_{j_t}}{\alpha \cdot p_{j_t} \cdot (s-1)} + \mathbb{E}_{t+1} \left[ (\tilde{a}_t^\top \tilde{\mu}_1 - \tilde{\alpha}_{\tilde{a}_t} \tilde{\mu}_1)^2 \right].
\]

Therefore, we have

\[
\mathbb{E}_{t}[G_c(I_t)] \leq \frac{\kappa_2}{s-1} + \kappa_2 \cdot \mathbb{E}_{\tilde{a}_t} \left[ \mathbb{E}_{t+1} \left[ (\tilde{a}_t^\top \tilde{\mu}_1 - \tilde{\alpha}_{\tilde{a}_t} \tilde{\mu}_1)^2 \right] \right] \tag{96}
\]

for some constant \( \kappa_2 > 0 \). Our proof is completed by combining (95) and (96). \(\square\)
Proof of Lemma 10. From (46), $S$ refers to the space spanned by $A$, which is a subspace of $\mathbb{R}^m$. We further denote by $S^\perp$ the orthogonal complement subspace of $S$ in the Euclidean space $\mathbb{R}^m$, i.e., $\mathbb{R}^m = S + S^\perp$ and $u^\top v = 0$ for any $u \in S$ and any $v \in S^\perp$.

Following basics in linear algebra, any vector $\mu \in \mathbb{R}^m$ can be decomposed uniquely as $\mu = \mu_S + \mu_{S^\perp}$, where $\mu_S \in S$ and $\mu_{S^\perp} \in S^\perp$. Then the projection of $\mu$ to the subspace $S$ can be given as $\mathcal{P}_S(\mu) = \mu_S$. Following this decomposition, for any $c \geq 0$, we have that

$$\min_{\mu \geq 0} L_{\text{c,t+1}}^{\text{Ftd}} = \min_{\mu_S + \mu_{S^\perp} \in S^\perp} \left( \frac{c}{T-t} \right)^\top \mu_S + \left( \frac{c}{T-t} \right)^\top \mu_{S^\perp} + \mathbb{E}(\tilde{r}, \tilde{a})[\tilde{r} - \tilde{a}^\top \mu_S]^+$$

where we note that $a^\top \mu_{S^\perp} = 0$ for any $a \in A$ and

$$h(\mu) = \min_{\mu \in S^\perp} \left( \frac{c}{T-t} \right)^\top \mu_{S^\perp} \quad \text{s.t.} \quad \mu_{S^\perp} + \mu_S \in \Omega.$$ 

Note that if the optimization problem that defines $h(\mu)$ is infeasible for some $\mu_S$, then we simply set $h(\mu_S) = +\infty$. Denote by $\hat{S} \subset S \cap \Omega$ the convex set such that $h(\mu)$ is finite for all $\mu \in \hat{S}$. Therefore, the dual problem $\min_{\mu \in S^\perp} L_{\text{c,t+1}}^{\text{Ftd}}$ can be transferred into a minimization problem over the set $\hat{S}$, i.e.,

$$\min_{\mu \in \hat{S}} L_{\text{c,t+1}}^{\text{Ftd}} = \min_{\mu \in \hat{S}} \hat{G}_{e,t+1}(\mu) := \left( \frac{c}{T-t} \right)^\top \mu + h(\mu) + \mathbb{E}(\tilde{r}, \tilde{a})[\tilde{r} - \tilde{a}^\top \mu]^+$$

Following the same way, the minimization of the sample average dual function $L_{\text{c,t+1}}^{\text{Off}}(\mu)$ in (42) can be formulated as

$$\min_{\mu \in \Omega} L_{\text{c,t+1}}^{\text{Off}} = \min_{\mu \in \hat{S}} \hat{G}_{e,t+1}(\mu) := \left( \frac{c}{T-t} \right)^\top \mu + h(\mu) + \frac{1}{T-t} \sum_{t=0}^{T} [\tilde{r}_t - \tilde{a}_t^\top \mu]^+.$$ 

We denote by

$$\mu^*(c') \in \arg \min_{\mu \in \hat{S}} \hat{G}_{c',t+1}(\mu) \quad \text{and} \quad \hat{\mu}(c') \in \arg \min_{\mu \in \hat{S}} \hat{G}_{c',t+1}(\mu)$$

for any $c' \geq 0$, where $\hat{\mu}(c')$ is a random variable whose value depends on $I_{t+1}$. Clearly, we have that

$$\mu^*(c) = \mathcal{P}_S(\hat{\mu}), \quad \text{and} \quad \hat{\mu}(c) = \mathcal{P}_S(\hat{\mu}).$$

Thus, in order to bound $\mathbb{E}[(a_i^\top \hat{\mu}_1 - \alpha_i^\top \hat{\mu}_2)^2]$, it is sufficient to bound

$$\mathbb{E}[(\mu^*(c') - \hat{\mu}(c'))^2]$$

for any $c' \geq 0$, where the $\mathbb{E}[\cdot]$ is taken over $I_{t+1}$. In what follows, we will consider bounding (99). Since we will assume that $c'$ is now fixed, for notation simplicity, we will drop $c'$ in the expression of $\mu^*(c')$, $\hat{\mu}(c')$, $\hat{G}_{c',t+1}(\mu)$, and $\hat{G}_{c',t+1}(\mu)$. We simply denote $\mu^*$, $\hat{\mu}$, $\hat{G}(\mu)$ and $\hat{G}_{t+1}(\mu)$.

We first note that the function $h(\mu)$ is a convex function over $\mu \in \hat{S}$. We also note that the function $\hat{G}$ is simply a re-formulation of the function $\hat{L}$ after projecting the decision variable onto the subspace $S$. Thus, Lemma 7 implies the following second-order growth condition for the function $\hat{G}$, i.e.,

$$\hat{G}(\mu) - \hat{G}(\mu^*) \geq \frac{\alpha \beta}{2} \|\mu - \mu^*\|_2^2$$

for any $\mu \in \hat{S} \subset S$. (100)
which implies that $\mu^*$ is unique. Thus, from Lemma 2.1 in Shapiro (1992), we know that
\[
\|\hat{\mu} - \mu^*\|_2 \leq \frac{2}{\alpha_2^2} \sup_{\mu \in \hat{S} \cap B(\mu^*, r_0)} \left\{ \frac{\bar{G}_{t+1}(\mu) - \tilde{G}(\mu) + \tilde{G}(\mu^*)}{\|\mu - \mu^*\|_2} \right\}
\]
where $r_0 = \|\hat{\mu} - \mu^*\|_2$ and $B(\mu^*, r_0) = \{ \mu : \|\mu - \mu^*\|_2 \leq r_0 \}$. Therefore, in order to further bound $\|\hat{\mu} - \mu^*\|_2$, it is sufficient to bound the term
\[
\sup_{\mu \in \hat{S} \cap B(\mu^*, r_0)} \left\{ \frac{1}{T-t} \sum_{t=1+1}^{T} [\tilde{r} - \tilde{a}^\top \mu]^* - \tilde{E}(\tilde{r}, \tilde{a})] [\tilde{r} - \tilde{a}^\top \mu]^* - \frac{1}{T-t} \sum_{t=1+1}^{T} [\tilde{r} - \tilde{a}^\top \mu^*]^* + \tilde{E}(\tilde{r}, \tilde{a})] [\tilde{r} - \tilde{a}^\top \mu^*]^* \right\}.
\]
We denote by the function
\[
g(\mu, (r, a)) := [r - a^\top \mu]^* + \alpha d^2 \cdot \|\mu - \mu^*\|_2^2.
\]
Moreover, we denote by
\[
\tilde{\delta}(\mu) = \frac{1}{T-t} \sum_{t=1+1}^{T} g(\mu, (\tilde{r}, \tilde{a}, \mu)) - \tilde{E}(\tilde{r}, \tilde{a})] [g(\mu, (\tilde{r}, \tilde{a}))]
\]
which is a random variable that depends on $I_{t+1}$. Then, the above upper bound over $\|\hat{\mu} - \mu^*\|_2$ can be re-formulated as
\[
\|\hat{\mu} - \mu^*\|_2 \leq \frac{2}{\alpha_2^2} \sup_{\mu \in \hat{S} \cap B(\mu^*, r_0)} \left\{ \frac{\tilde{\delta}(\mu) - \tilde{\delta}(\mu^*)}{\|\mu - \mu^*\|_2} \right\}.
\]
We denote by $\hat{E}$ the set of $\mu$ such that $\tilde{\delta}(\mu)$ is non-differentiable. For simplicity, we denote by $\hat{B}(\mu^*, r_0) = \hat{S} \cap B(\mu^*, r_0)$. Note that $\tilde{\delta}(\mu)$ is Liptschitz continuous. Then, from the mean value theorem for Lipstchitz function (Clarke (1990) p41), we have that
\[
\|\hat{\mu} - \mu^*\|_2 \leq \frac{4}{\alpha_2^2 \beta^2} \sup_{\mu \in \hat{B}(\mu^*, r_0) \setminus \hat{E}} \{ \|\nabla \tilde{\delta}(\mu)\|_2^2 \} = \frac{4}{\alpha_2^2 \beta^2} \cdot (\|\nabla \tilde{\delta}(\mu^*)\|_2^2 + \sup_{\mu \in \hat{B}(\mu^*, r_0) \setminus \hat{E}} \{ \|\nabla \tilde{\delta}(\mu) - \nabla \tilde{\delta}(\mu^*)\|_2^2 \})
\]
We have the following claim over the term $\|\nabla \tilde{\delta}(\mu^*)\|_2^2$, which is proved at the end of this proof.

**Claim 5.** It holds that
\[
\mathbb{E}[\|\nabla \tilde{\delta}(\mu^*)\|_2^2] \leq \frac{\bar{d}^2}{T-t}
\]
where the expectation $\mathbb{E}[\cdot]$ takes over $I_{t+1} = \{(\tilde{r}_{t+1}, \tilde{a}_{t+1}), \ldots, (\tilde{r}_T, \tilde{a}_T)\}$ that defines $\tilde{\delta}(\cdot)$ in (101).

Regarding the term $\sup_{\mu \in \hat{B}(\mu^*, r_0) \setminus \hat{E}} \{ \|\nabla \tilde{\delta}(\mu) - \nabla \tilde{\delta}(\mu^*)\|_2^2 \}$, following the steps in Shapiro (1993), we can show that
\[
\sup_{\mu \in \hat{B}(\mu^*, r_0) \setminus \hat{E}} \{ \|\nabla \tilde{\delta}(\mu) - \nabla \tilde{\delta}(\mu^*)\|_2 \} \to 0 \text{ in probability as } T - t \to \infty.
\]
As a result, there exists a constant $t_0$ such that as long as $T - t \geq t_0$, we have that
\[
\sup_{\mu \in \hat{B}(\mu^*, r_0) \setminus \hat{E}} \{ \|\nabla \tilde{\delta}(\mu) - \nabla \tilde{\delta}(\mu^*)\|_2^2 \} \leq \frac{1}{2} \cdot (\frac{1}{T-t} + r_0^2)
\]
with a high probability. We formalize the above step in the following claim for completeness, by making explicit the constant term $t_0$, as well as the “high probability”.

**Claim 6.** It holds that
\[
P \left( \sup_{\|\mu - \mu^*\|_2 \leq r_0} \frac{\|\nabla \tilde{\delta}(\mu) - \nabla \tilde{\delta}(\mu^*)\|_2}{\sqrt{1 + 3\alpha d^2 \cdot r_0}} \geq \frac{\alpha \beta}{12 \alpha d^2} \right) \leq \delta(\tilde{\alpha}, \tilde{d}, m, \gamma) \cdot \frac{\log(T-t)}{(T-t)^{\frac{1}{3}}}
\]
as long as $T - t \geq t_1(\tilde{\alpha}, \tilde{d})$, where $\delta_1(\tilde{\alpha}, \tilde{d}, m, \gamma)$ is a constant that depends solely on the parameters $\tilde{\alpha}, \tilde{d}, m, \gamma$ and $t_1(\tilde{\alpha}, \tilde{d})$ is a constant that depends solely on the parameters $\tilde{\alpha}, \tilde{d}$. 

The proof of Claim 6 is relegated to the end of this proof. Denote by
\[
\tilde{z} = \sup_{\|\mu - \mu^*\|_2 \leq r_0} \frac{\|\nabla \hat{\delta} (\mu) - \nabla \hat{\delta} (\mu^*)\|_2}{\sqrt{\|T\|} + 3\tilde{\alpha}d^2 \cdot r_0}.
\]
We also denote by three events
\[
E_1 = \{\tilde{z} \leq \frac{\alpha \beta}{12 \tilde{\alpha}d^2}\} \quad \text{and} \quad E_2 = \{\frac{\alpha \beta}{12 \tilde{\alpha}d^2} < \tilde{z} \quad \text{and} \quad r_0 \leq \frac{1}{(T - t)^{\frac{1}{20}}}\} \quad \text{and} \quad E_3 = \{\frac{\alpha \beta}{12 \tilde{\alpha}d^2} < \tilde{z} \quad \text{and} \quad r_0 > \frac{1}{(T - t)^{\frac{1}{20}}}\}
\]
Clearly, we have
\[
\mathbb{E}[r_0^2] = P(E_1) \cdot \mathbb{E}[r_0^2 | E_1] + P(E_2) \cdot \mathbb{E}[r_0^2 | E_2] + P(E_3) \cdot \mathbb{E}[r_0^2 | E_3]
\]
Now, from (102), Claim 5, and the definition of the event \(E_1\), we have
\[
P(E_1) \cdot \mathbb{E}[r_0^2 | E_1] \leq \frac{4}{\alpha^2 \beta^2} \cdot P(E_1) \cdot \mathbb{E}[\|\nabla \hat{\delta} (\mu^*)\|_2^2 | E_1] + P(E_1) \cdot \frac{1}{36 \tilde{\alpha}^2 d^4} \cdot \mathbb{E} \left[ \left( \frac{1}{\sqrt{\|T\|}} + 3\tilde{\alpha}d^2 \cdot r_0 \right)^2 | E_1 \right] \leq \frac{4}{\alpha^2 \beta^2} \cdot \mathbb{E}[\|\nabla \hat{\delta} (\mu^*)\|_2^2] + \frac{1}{2} \cdot P(E_1) \cdot \mathbb{E}[r_0^2 | E_1] + \frac{1}{18 \tilde{\alpha}^2 d^4 \cdot (T - t)}
\]
which implies that
\[
P(E_1) \cdot \mathbb{E}[r_0^2 | E_1] \leq \frac{8d^2}{\alpha^2 \beta^2 \cdot (T - t)} + \frac{1}{9 \tilde{\alpha}^2 d^4 \cdot (T - t)}
\]
For the term \(P(E_2) \cdot \mathbb{E}[r_0^2 | E_2]\), from the definition of the event \(E_2\) and Claim 6, we have
\[
P(E_2) \cdot \mathbb{E}[r_0^2 | E_2] \leq P(E_2) \cdot \frac{1}{(T - t)^{\frac{1}{20}}} \leq \delta(\tilde{\alpha}, \tilde{\alpha}, m, \gamma) \cdot \frac{\log(T - t)}{(T - t)^{\frac{1}{20}}} \cdot \frac{1}{(T - t)^{\frac{1}{20}}} \leq \frac{1}{(T - t)^{\frac{1}{20}}}
\]
as long as \(T - t > t_2(\tilde{\alpha}, \tilde{\alpha}, m, \gamma)\), where \(t_2(\tilde{\alpha}, \tilde{\alpha}, m, \gamma)\) is a constant that depends solely on parameters \(\tilde{\alpha}, \tilde{\alpha}, m, \gamma\).

We want to bound the probability that event \(E_3\) happens. Note that we have
\[
\frac{\alpha \beta}{2} \cdot \|\mu - \mu^*\|_2^2 \leq \frac{1}{T - t} \sum_{r = t + 1}^T \left[ \hat{r}^r - \tilde{\alpha}^\top \mu^* \right] + \mathbb{E}[r, \tilde{a}][\hat{r}^r - \tilde{\alpha}^\top \mu^*] - \frac{1}{T - t} \sum_{r = t + 1}^T \left[ \hat{r}^r - \tilde{\alpha}^\top \mu^* \right] + \mathbb{E}[r, \tilde{a}][\hat{r}^r - \tilde{\alpha}^\top \mu^*].
\]
By modifying the approach in Li and Ye (2021), we have the following lemma, which follows Proposition 2 and Proposition 3 of Li and Ye (2021) and is proved at the end of this proof.

**Lemma 16.** It holds that
\[
P \left( \|\mu^\star - \mu^*\|_2^2 \geq \epsilon^2 \right) \leq K \cdot \log \left( \frac{1}{\epsilon} \right) \exp(-\epsilon^2 (T - t)).
\]
for a constant \(K\) that only depends on parameters \(\alpha, \beta, \tilde{\alpha}, \tilde{\alpha}\)

By setting \(\epsilon = \frac{1}{(T - t)^{\frac{1}{20}}}\) in Lemma 16, we have that
\[
P(E_3) \leq P \left( \|\mu^\star - \mu^*\|_2^2 \geq \left( \frac{1}{(T - t)^{\frac{1}{20}}} \right)^2 \right) \leq K' \cdot \log(T - t) \cdot \exp(- (T - t)^{\frac{1}{40}}) \leq \frac{K''}{T - t}
\]
as long as \(T - t > t_3\), where \(K', K'', t_3\) are constants. Therefore, combining (104), (106), (107), and (109), we have
\[
\mathbb{E}[\|\mu^\star - \mu^*\|_2^2] = \mathbb{E}[r_0^2] \leq \frac{\delta_2}{T - t}
\]
as long as \(T - t > t_4\), where \(\delta_2\) and \(t_4\) are constants. Therefore, our proof is completed. \(\square\)
Proof of Claim 5. Note that
\[ \nabla \delta(\mu^*) = -\frac{1}{T-t} \cdot \sum_{\tau=t+1}^{T} \bar{a}_\tau \cdot \mathbb{I}_{(\bar{r}_\tau \geq \bar{a}_\tau \mu^*)} + \mathbb{E}_a[\bar{a} \cdot (1 - F(\bar{a}^T \mu^* | \bar{a}))] \]
Clearly, for each \( \tau = t+1, \ldots, T \), we have
\[ \mathbb{E}_{(\bar{r}, \bar{a}_\tau)}[\bar{a}_\tau \cdot \mathbb{I}_{(\bar{r}_\tau \geq \bar{a}_\tau \mu^*)}] = \mathbb{E}_a[\bar{a} \cdot (1 - F(\bar{a}^T \mu^* | \bar{a}))]. \]
Thus, we have
\[ \mathbb{E}[\|\nabla \delta(\mu^*)\|_2^2] = \frac{1}{T-t} \cdot \text{Var}(\bar{a}_\tau \cdot \mathbb{I}_{(\bar{r}_\tau \geq \bar{a}_\tau \mu^*)}) \leq \frac{d^2}{T-t} \]
which completes our proof. \( \square \)

Proof of Claim 6. We denote by
\[ \psi((r, a), \mu) = \nabla g(\mu, (r, a)) - \nabla g(\mu^*, (r, a)) \]
We also denote by
\[ \lambda(\mu) = \mathbb{E}_{(r, a)}[\psi((r, a), \mu)] \]
and
\[ u((r, a), \mu, d) = \sup_{\|\mu - \mu^*\|_2 \leq d} \|\psi((r, a), \mu) - \psi((r, a), \mu^*)\|_2 \]
Clearly, we have that
\[ \lambda(\mu^*) = 0 \] (110)
Thus, condition (N-2) in Huber (1967) is satisfied. Also, for any \( \mu \), from definition of the functions \( \psi \) and \( \lambda \), we have that
\[ \|\lambda(\mu)\|_2 = \|2\bar{a}d^2(\mu - \mu^*) + \mathbb{E}_a[\bar{a} \cdot (F(\bar{a}^T \mu | \bar{a}) - F(\bar{a}^T \mu^* | \bar{a}))]\|_2 \]
Note that
\[ F(\bar{a}^T \mu | \bar{a}) - F(\bar{a}^T \mu^* | \bar{a}) \leq \bar{a} \cdot |\bar{a}^T \mu - \bar{a}^T \mu^*| \leq \bar{a}d \cdot \|\mu - \mu^*\|_2 \]
Thus, we have
\[ \|\lambda(\mu)\|_2 \geq 2\bar{a}d^2 \cdot \|\mu - \mu^*\|_2 - \|\mathbb{E}_a[\bar{a} \cdot (F(\bar{a}^T \mu | \bar{a}) - F(\bar{a}^T \mu^* | \bar{a}))]\|_2 \]
\[ \geq 2\bar{a}d^2 \cdot \|\mu - \mu^*\|_2 - \bar{a}d^2 \cdot \|\mu - \mu^*\|_2 = \bar{a}d^2 \cdot \|\mu - \mu^*\|_2 \] (111)
for any \( d \geq 0 \) and \( \|\mu - \mu^*\|_2 \leq d \). Thus, we verify that condition (N-3) (i) in Huber (1967) holds. Moreover, we have
\[ \|\psi((r, a), \mu) - \psi((r, a), \mu^*)\|_2 = \|a \cdot \mathbb{I}_{(\bar{a} \geq \bar{a}^T \mu^*)} - a \cdot \mathbb{I}_{(\bar{a}^T \mu \leq \bar{a}^T \mu^*)}\| \leq \|a \cdot (\mathbb{I}_{(a \geq \bar{a}^T \mu)} + \mathbb{I}_{(a \leq \bar{a}^T \mu^*)})\| \leq \bar{a} \]
which implies that
\[ \mathbb{E}_{(\bar{r}, \bar{a})}[u((\bar{r}, \bar{a}), \mu, d)] \leq \mathbb{E}_a[\|\bar{a}\|_2 \cdot 2\bar{a}d] \leq 2\bar{a}d^2 \cdot d \] (113)
Similarly, we have
\[ \mathbb{E}_{(\bar{r}, \bar{a})}[u((\bar{r}, \bar{a}), \mu, d)] \leq \mathbb{E}_a[\|\bar{a}\|_2^3 \cdot 2\bar{a}d] \leq 2\bar{a}d^3 \cdot d \]
Thus, we verify conditions (N-3) (ii) and (iii) in Huber (1967) for any \( d_0 \geq 0 \). Clearly, \( E_{(r, \bar{a})} [\psi((\bar{r}, \bar{a}), \mu^*)] \) is finite, which verifies condition (N-4) in Huber (1967). Thus, denote by

\[
Z_{T-t}(\mu, \mu^*) = \frac{1}{T-t} \sum_{t=r+1}^{T} (\psi((\bar{r}, \bar{a}, r), \mu) - \psi((\bar{r}, \bar{a}, r), \mu^*)) - \lambda(\mu) + \lambda(\mu^*)
\]

By Lemma 15, for any \( \epsilon > 0 \) and any \( \gamma' \in \left( \frac{1}{2}, 1 \right) \), we have

\[
P \left( \sup_{\| \mu - \mu^* \| \leq d_0} Z_{T-t}(\mu, \mu^*) \geq 2\epsilon \right) \leq 2\bar{a}d^2 \epsilon^{-2}(T-t)^{-\gamma'}
\]

\[
+ (T-t)^{\gamma'-1} \left( \frac{2\alpha \bar{a} d}{4\alpha^2 d^2 q(1-q)} + \frac{2\alpha \bar{a} d}{36\alpha^2 d^4 q^2(1-q)^2} \right) \left( \gamma' \left( \frac{\log(T-t) + \log d_0}{\log(1-q)} \right) + 1 \right) (2M)^m
\]

as long as \( T-t \geq n_0 \), where \( n_0 \) satisfying \( n_0^{\gamma'-\frac{1}{2}} = \frac{4\alpha \bar{a} d}{\epsilon} \) and \( \gamma' \in \left( \frac{1}{2}, 1 \right) \) is an arbitrary number. Moreover, we set \( M \geq (6\bar{a}d^2)/(\epsilon \bar{a}d^2) = 6\bar{a}/\epsilon \) and \( q = 1/M \).

We first specify \( \epsilon = \frac{\alpha \bar{a} d}{24\alpha \bar{a} d} \) and \( \gamma' = 13/25 \). Then, (114) will become

\[
P \left( \sup_{\| \mu - \mu^* \| \leq d_0} Z_{T-t}(\mu, \mu^*) \geq \frac{\alpha \bar{a} d}{12\alpha \bar{a} d} \right) \leq 2\bar{a}d^2 \left( \frac{12\bar{a}d^2}{\alpha \bar{a} d} \right)^2 (T-t)^{-\frac{13}{25}}
\]

\[
+ (T-t)^{\frac{13}{25}} \left( \frac{2\alpha \bar{a} d}{4\alpha^2 d^2 q(1-q)} + \frac{2\alpha \bar{a} d}{36\alpha^2 d^4 q^2(1-q)^2} \right) \left( \frac{\log(T-t) + \log d_0}{25 \log(1-q)} \right) + 1 \right) (2M)^m
\]

as long as \( T-t \geq (96\bar{a}d^2/(\alpha \bar{a} d))^b_{00} \), \( M = 144 \bar{a}d^2/(\alpha \bar{a} d) \) and \( q = 1/M = \alpha \bar{a} d/(144 \bar{a}d^2) \). Denote by \( \delta(\bar{a}, \bar{d}, m) \) a constant that depends solely on \( \bar{a}, \bar{d} \) and \( m \), i.e.,

\[
\delta(\bar{a}, \bar{d}, m) = 2\bar{a}d^3 \left( \frac{12\bar{a}d^2}{\alpha \bar{a} d} \right)^2 + \left( \frac{2\alpha \bar{a} d}{4\alpha^2 d^2 q(1-q)} + \frac{2\alpha \bar{a} d}{36\alpha^2 d^4 q^2(1-q)^2} \right) \left( \frac{13}{25 \log(1-q)} \right) + 1 \right) (2M)^m
\]

with \( M = 144 \bar{a}d^2/(\alpha \bar{a} d) \) and \( q = 1/M = \alpha \bar{a} d/(144 \bar{a}d^2) \). Then, from (115), we have

\[
P \left( \sup_{\| \mu - \mu^* \| \leq d_0} Z_{T-t}(\mu, \mu^*) \geq \frac{\alpha \bar{a} d}{12\alpha \bar{a} d} \right) \leq \delta(\bar{a}, \bar{d}, m) \cdot \frac{\log(T-t) + \log d_0}{(T-t)^{\frac{13}{25}}}
\]

as long as \( T-t \geq (96\bar{a}d^2/(\alpha \bar{a} d))^b_{00} \). Note that

\[
\nabla \delta(\mu) - \nabla \delta(\mu^*) = \frac{1}{T-t} \cdot \sum_{t=r+1}^{T} (\psi((\bar{r}, \bar{a}, r), \mu) - \psi((\bar{r}, \bar{a}, r), \mu^*)) - \lambda(\mu) + \lambda(\mu^*)
\]

We know that (117) implies that

\[
P \left( \sup_{\| \mu - \mu^* \| \leq d_0} \left\| \nabla \delta(\mu) - \nabla \delta(\mu^*) \right\| \right) \leq \frac{\alpha \bar{a} d}{12\alpha \bar{a} d} \cdot \frac{\log(T-t) + \log d_0}{(T-t)^{\frac{13}{25}}}
\]

From Assumption 1, we have \( \| \mu \|_2 \leq \gamma \sqrt{m} \) and \( \| \mu^* \|_2 \leq \gamma \sqrt{m} \), which implies that \( r_0 = \| \mu - \mu^* \|_2 \leq 2\gamma \sqrt{m} \).

Therefore, we can simply set \( d_0 = 2\gamma \sqrt{m} \) and we know that \( r_0 \) is always upper bounded by \( d_0 \), which implies that

\[
P \left( \sup_{\| \mu - \mu^* \| \leq r_0} \left\| \nabla \delta(\mu) - \nabla \delta(\mu^*) \right\| \right) \leq \frac{\alpha \bar{a} d}{12\alpha \bar{a} d} \cdot \frac{\log(T-t) + \log(2\gamma \sqrt{m})}{(T-t)^{\frac{13}{25}}}
\]

where the first inequality follows from \( r_0 \leq d_0 = 2\gamma \sqrt{m} \) almost surely and the second inequality follows from (118) by setting \( d_0 = 2\gamma \sqrt{m} \). Therefore, we denote by \( \delta(\bar{a}, \bar{d}, m, \gamma) = \delta(\bar{a}, \bar{d}, m) \cdot \log(2\gamma \sqrt{m}) \) and our proof of (103) is completed by noting that

\[
\| \lambda(\mu) \|_2 = \| 2\bar{a}d^2 (\mu - \mu^*) + E_{(r, \bar{a})} [\bar{a} \cdot (F(\bar{a}^T \mu) \bar{a}) - F(\bar{a}^T \mu^*) \bar{a})] \|_2 \leq 3\bar{a}d^2 \cdot \| \mu - \mu^* \|_2.
\]

□
Proof of Lemma 16. Following Li and Ye (2021), we denote by
\[ \phi((r, a), \mu) = -a^\top \mathbb{1}(r > a^\top \mu). \]

Then, from Lemma 1 of Li and Ye (2021), we have
\[
\mathbb{E}(\hat{r}, \hat{a})[\bar{r} - \bar{\mu}^\top \mu]^+ - \mathbb{E}(\hat{r}, \hat{a})[\bar{r} - \bar{\mu}^\top \mu]^+ = \mathbb{E}(\hat{r}, \hat{a})[\phi((\hat{r}, \hat{a}), \mu^*)] \cdot (\hat{\mu} - \mu^*) + \mathbb{E}(\hat{r}, \hat{a}) \left[ \int_{\bar{\mu}}^{\hat{\mu}^*} (\mathbb{1}(r > v) - \mathbb{1}(\hat{r} > \hat{\mu}^\top \mu^*)) dv \right].
\]

Also, from Lemma 2 of Li and Ye (2021), we have
\[
\frac{1}{T-t} \sum_{t=t+1}^{T} \phi((\hat{r}, \hat{a}), \mu^*) \cdot (\hat{\mu} - \mu^*) = \frac{1}{T-t} \sum_{t=t+1}^{T} \phi((\hat{r}, \hat{a}), \mu^*) \cdot (\hat{\mu} - \mu^*) + \frac{1}{T-t} \sum_{t=t+1}^{T} \int_{\hat{a}^\top \hat{\mu}}^{\hat{\mu}^*} (\mathbb{1}(\hat{r} > v) - \mathbb{1}(\hat{r} > \hat{a}^\top \mu^*)) dv
\]

Plugging (120) and (121) into (108), we have that
\[
\frac{a_B^2}{2} \| \hat{\mu} - \mu^* \|^2 \leq \frac{1}{T-t} \sum_{t=t+1}^{T} \phi((\hat{r}, \hat{a}), \mu^*) \cdot (\hat{\mu} - \mu^*) \cdot \sum_{t=t+1}^{T} \phi((\hat{r}, \hat{a}), \mu^*) - \mathbb{E}(\hat{r}, \hat{a})[\phi((\hat{r}, \hat{a}), \mu^*)] \cdot (\hat{\mu} - \mu^*) + \frac{1}{T-t} \sum_{t=t+1}^{T} \int_{\hat{a}^\top \hat{\mu}}^{\hat{\mu}^*} (\mathbb{1}(\hat{r} > v) - \mathbb{1}(\hat{r} > \hat{a}^\top \mu^*)) dv - \mathbb{E}(\hat{r}, \hat{a}) \left[ \int_{\hat{a}^\top \hat{\mu}}^{\hat{\mu}^*} (\mathbb{1}(\hat{r} > v) - \mathbb{1}(\hat{r} > \hat{a}^\top \mu^*)) dv \right]
\]

We proceed by bounding the term I and term II separately.

**Bound I:** by Matrix Hoeffding’s inequality, we have that
\[
P \left( \frac{1}{T-t} \sum_{t=t+1}^{T} \phi((\hat{r}, \hat{a}), \mu^*) - \mathbb{E}(\hat{r}, \hat{a})[\phi((\hat{r}, \hat{a}), \mu^*)] \right) \leq m \cdot \exp \left( -\frac{\epsilon^2 \cdot (T-t)}{d^2} \right)
\]

which implies that
\[
P (|I| \geq \epsilon \cdot \| \hat{\mu} - \mu^* \|) \leq m \cdot \exp \left( -\frac{\epsilon^2 \cdot (T-t)}{d^2} \right)
\]

**Bound II:** We define a function
\[
\eta((r, a), \mu_1, \mu_2) = \int_{\alpha^\top \mu_2}^{\alpha^\top \mu_1} (\mathbb{1}(r > v) - \mathbb{1}(r > \alpha^\top \mu^*)) dv
\]

for any \( \mu_1, \mu_2 \in \mathcal{S} \). We have
\[
II = \frac{1}{T-t} \sum_{t=t+1}^{T} \eta((\hat{r}, \hat{a}), \hat{\mu}, \mu^*) - \mathbb{E}(\hat{r}, \hat{a})[\eta((\hat{r}, \hat{a}), \hat{\mu}, \mu^*)]
\]

We utilize a splitting scheme to split the set \( \{ \mu : \| \mu - \mu^* \| \leq \sqrt{m} \cdot \gamma \} \) into disjoint cubes, similar to Huber (1967) and Li and Ye (2021). It is clear to see that \( \| \hat{\mu} - \mu^* \| \leq \sqrt{m} \cdot \gamma \). The idea is to divide the region \( \{ \mu : \| \mu - \mu^* \| \leq \sqrt{m} \cdot \gamma \} \) into a slowly increasing number of smaller cubes. We consider the concentric cubes:
\[
C_k = \{ \mu : \| \mu - \mu^* \| \leq (1-q)^k \cdot \sqrt{m} \gamma \}, \ k = 1, \ldots, k_0
\]
where \( q \) is a constant such that \( \bar{\alpha} \cdot \bar{d}^2 \cdot \frac{q(1+q)}{(1-q)^2} \leq \frac{\alpha^2}{32}, \) and \( k_0 \) is a constants such that \( (1-q)^k \cdot \sqrt{m\gamma} = \epsilon. \) We then further divide the region \( C_{-1} \setminus C_k \) into cubes \( \{\Omega_{\alpha}\}^k_{l=1} \) with edges of length \( 2(1-q)^k \cdot q \cdot \sqrt{m\gamma} \) such that the centers of these cubes \( \xi_{l\ell} \) satisfies

\[
\|\xi_{l\ell}\| = (1-q)^{k-1}(1-q)\frac{\sqrt{m\gamma}}{2}, \quad l = 1, \ldots, k_l.
\]

In total, there are no more than \( k_0 \cdot \left(\frac{q}{2}\right)^m \) number of cubes. We now denote by \( \bar{\Omega} \) the cube that contains \( \tilde{\mu} \) and denote by \( \tilde{\xi} \) the center of the cube \( \bar{\Omega}. \) Note that both \( \bar{\Omega} \) and \( \tilde{\xi} \) are random, where the randomness arises from \( \tilde{\mu}. \)

\[
\frac{1}{T-t} \cdot \sum_{r=1}^{T} \eta((\tilde{r}_r, \tilde{a}_r), \tilde{\mu}, \mu^*) - \mathbb{E}_{(\tilde{r}, \tilde{a})}[\eta((\tilde{r}, \tilde{a}), \tilde{\mu}, \mu^*)] = \frac{1}{T-t} \cdot \sum_{r=1}^{T} \eta((\tilde{r}_r, \tilde{a}_r), \tilde{\mu}, \tilde{\xi}) - \mathbb{E}_{(\tilde{r}, \tilde{a})}[\eta((\tilde{r}, \tilde{a}), \tilde{\mu}, \tilde{\xi})]
\]

+ \frac{1}{T-t} \cdot \sum_{r=1}^{T} \eta((\tilde{r}_r, \tilde{a}_r), \tilde{\xi}, \mu^*) - \mathbb{E}_{(\tilde{r}, \tilde{a})}[\eta((\tilde{r}, \tilde{a}), \tilde{\xi}, \mu^*)]

which implies that

\[
P \left( \frac{1}{T-t} \cdot \sum_{r=1}^{T} \eta((\tilde{r}_r, \tilde{a}_r), \tilde{\mu}, \mu^*) - \mathbb{E}_{(\tilde{r}, \tilde{a})}[\eta((\tilde{r}, \tilde{a}), \tilde{\mu}, \mu^*)] \geq \epsilon^2 + \epsilon \|\tilde{\mu} - \mu^*\|_2 + \frac{\alpha \beta \|\tilde{\mu} - \mu^*\|_2^3}{8} \right)
\]

\[
\leq P \left( \frac{1}{T-t} \cdot \sum_{r=1}^{T} \eta((\tilde{r}_r, \tilde{a}_r), \tilde{\mu}, \tilde{\xi}) - \mathbb{E}_{(\tilde{r}, \tilde{a})}[\eta((\tilde{r}, \tilde{a}), \tilde{\mu}, \tilde{\xi})] \geq \epsilon^2 + \epsilon \|\tilde{\mu} - \mu^*\|_2 + \frac{\alpha \beta \|\tilde{\mu} - \mu^*\|_2^3}{16} \right)
\]

+ \frac{1}{T-t} \cdot \sum_{r=1}^{T} \eta((\tilde{r}_r, \tilde{a}_r), \tilde{\xi}, \mu^*) - \mathbb{E}_{(\tilde{r}, \tilde{a})}[\eta((\tilde{r}, \tilde{a}), \tilde{\xi}, \mu^*)] \geq \epsilon^2 + \epsilon \|\tilde{\mu} - \mu^*\|_2 + \frac{\alpha \beta \|\tilde{\mu} - \mu^*\|_2^3}{16} \right)
\]

We have the following result.

**Claim 7.** It holds that

\[
P \left( \frac{1}{T-t} \cdot \sum_{r=1}^{T} \eta((\tilde{r}_r, \tilde{a}_r), \tilde{\mu}, \tilde{\xi}) - \mathbb{E}_{(\tilde{r}, \tilde{a})}[\eta((\tilde{r}, \tilde{a}), \tilde{\mu}, \tilde{\xi})] \geq \epsilon^2 + \epsilon \|\tilde{\mu} - \mu^*\|_2 + \frac{\alpha \beta \|\tilde{\mu} - \mu^*\|_2^3}{16} \right)
\]

\[
\leq k_0 \cdot \left(\frac{2}{q}\right)^m \cdot \exp \left( -\frac{\epsilon^2(T-t)q^2}{2d^2} \right)
\]

where \( q \) is a constant such that \( \bar{\alpha} \cdot \bar{d}^2 \cdot \frac{q(1+q)}{(1-q)^2} \leq \frac{\alpha^2}{32}, \) and \( k_0 \) is a constants such that \( (1-q)^k \cdot \sqrt{m\gamma} = \epsilon. \)

We also have the following result.

**Claim 8.** It holds that

\[
P \left( \frac{1}{T-t} \cdot \sum_{r=1}^{T} \eta((\tilde{r}_r, \tilde{a}_r), \tilde{\mu}, \mu^*) - \mathbb{E}_{(\tilde{r}, \tilde{a})}[\eta((\tilde{r}, \tilde{a}), \tilde{\mu}, \mu^*)] \geq \epsilon^2 + \epsilon \|\tilde{\mu} - \mu^*\|_2 + \frac{\alpha \beta \|\tilde{\mu} - \mu^*\|_2^3}{16} \right)
\]

\[
\leq k_0 \cdot \left(\frac{2}{q}\right)^m \cdot \exp \left( -\frac{\epsilon^2(T-t)q^2}{2d^2} \right)
\]

where \( q \) is a constant such that \( \bar{\alpha} \cdot \bar{d}^2 \cdot \frac{q(1+q)}{(1-q)^2} \leq \frac{\alpha^2}{32}, \) and \( k_0 \) is a constants such that \( (1-q)^k \cdot \sqrt{m\gamma} = \epsilon. \)

Combining Claim 7, Claim 8 and (125), we have that

\[
P \left( \|I\|^2 \geq \epsilon^2 + \epsilon \|\tilde{\mu} - \mu^*\|_2 + \frac{\alpha \beta \|\tilde{\mu} - \mu^*\|_2^3}{8} \right)
\]

\[
= P \left( \frac{1}{T-t} \cdot \sum_{r=1}^{T} \eta((\tilde{r}_r, \tilde{a}_r), \tilde{\mu}, \mu^*) - \mathbb{E}_{(\tilde{r}, \tilde{a})}[\eta((\tilde{r}, \tilde{a}), \tilde{\mu}, \mu^*)] \geq \epsilon^2 + \epsilon \|\tilde{\mu} - \mu^*\|_2 + \frac{\alpha \beta \|\tilde{\mu} - \mu^*\|_2^3}{8} \right)
\]

\[
\leq k_0 \cdot \left(\frac{2}{q}\right)^m \cdot \exp \left( -\frac{\epsilon^2(T-t)q^2}{2d^2} \right) + k_0 \cdot \left(\frac{2}{q}\right)^m \cdot \exp \left( -\frac{\epsilon^2(T-t)q^2}{2d^2} \right)
\]
where $q$ is a constant such that $\bar{\alpha} \cdot \bar{d}^2 \cdot \frac{q(1 + q)}{(1 - q)^2} \leq \frac{\alpha \beta}{2}$, and $k_0$ is a constants such that $(1 - q)^{k_0} \cdot \sqrt{m \gamma} = \epsilon$. Our bound of the term II is now completed.

We now use the bound on I and II to prove the probability bound on the event $\|\hat{\mu} - \mu^*\|^2 \geq \epsilon^2$. Combining (123) and (126), we have

$$P \left( I + II \geq 2\epsilon^2 + 2\epsilon \|\hat{\mu} - \mu^*\|_2 + \frac{\alpha \beta}{4} \|\hat{\mu} - \mu^*\|^2 \right)$$

where $c_2$ is a constant. Further note that

$$\frac{\alpha \beta}{2} \|\hat{\mu} - \mu^*\|^2 \leq I + II,$$

and

$$\frac{\alpha \beta}{2} \cdot \|\hat{\mu} - \mu^*\|^2 \geq 2\epsilon^2 + 2\epsilon \|\hat{\mu} - \mu^*\|_2 + \frac{\alpha \beta}{4} \|\hat{\mu} - \mu^*\|^2$$

implies that

$$\|\hat{\mu} - \mu^*\|^2 \geq \epsilon^2 \cdot \frac{4(2 + \sqrt{2\alpha \beta + 4})}{\alpha \beta}.$$

Therefore, we have that

$$P \left( \|\hat{\mu} - \mu^*\|^2 \geq \epsilon^2 \cdot \frac{4(2 + \sqrt{2\alpha \beta + 4})}{\alpha \beta} \right) \leq m \cdot \exp \left( -\epsilon^2 \cdot \frac{(T - t)}{d^2} \right)$$

By substituting $\epsilon' = \epsilon \cdot \frac{2\sqrt{2\alpha \beta + 4}}{\sqrt{\alpha \beta}}$ and noting that $k_0 = O(\log(\frac{1}{\epsilon}))$, we have that

$$P \left( \|\hat{\mu} - \mu^*\|^2 \geq (\epsilon')^2 \right) \leq K \cdot \log \left( \frac{1}{\epsilon'} \right) \exp \left( - (\epsilon')^2 (T - t) \right)$$

for a constant $K$ that only depends on parameters $\alpha, \beta, \bar{\alpha}, \bar{d}$, which completes our proof. \hfill \Box

**Proof of Claim 7.** We denote by

$$g((r, a), \mu_1, d) = \sup_{\mu_2: \|\mu_2 - \mu_1\|_2 \leq d} \eta((r, a), \mu_1, \mu_2).$$

Then we have

$$\frac{1}{T - t} \cdot \sum_{\tau = \tau + 1}^{T} \eta((\hat{r}, \hat{a}, \hat{\mu}, \hat{\xi})) - \mathbb{E}_{(\hat{r}, \hat{a})} [\eta((\hat{r}, \hat{a}, \hat{\mu}, \hat{\xi}))] \leq \frac{1}{T - t} \cdot \sum_{\tau = \tau + 1}^{T} g((\hat{r}, \hat{a}, \hat{\xi}, \hat{d})) + \mathbb{E}_{(\hat{r}, \hat{a})} [g((\hat{r}, \hat{a}, \hat{\xi}, \hat{d}))]$$

where $\hat{d}$ denotes the edge length of the cube $\hat{\Omega}$. We now consider the cubes $\{\Omega_\delta\}$ and the center concentric cubes $C_\delta$ separately.

(i). If the cube $\Omega \in \{\Omega_\delta\}$, then, we note that

$$\mathbb{E}_{(\hat{r}, \hat{a})} [g((\hat{r}, \hat{a}, \hat{\xi}, \hat{d}))] \leq \mathbb{E}_{\hat{a}} \left[ \sup_{\mu_2: \|\mu_2 - \mu_1\|_2 \leq \hat{d}} \int_{\hat{a}}^{\hat{a} + \hat{\xi}} \hat{\xi} \cdot \hat{\mu} \cdot d\hat{v} \int_{\hat{a}}^{\hat{a} + \hat{\xi}} (\mathbb{1}\{r > v\} - \mathbb{1}\{r > \hat{a} + \hat{\xi}\})dF(r) \hat{a} \right] \leq \hat{a} \hat{d}^2 \cdot \hat{d} \cdot (\|\mu^* - \hat{\xi}\|_2 + \hat{d})$$
By definition, we have $\bar{d} = q \cdot \|\mu^* - \bar{\xi}\|_2$ and $\|\mu^* - \bar{\mu}\|_2 \geq \|\mu^* - \bar{\xi}\|_2 - \bar{d}$. Then, we have
\[
\mathbb{E}_{(\bar{r}, \bar{a})}[g((\bar{r}, \bar{a}), \bar{\xi}, \bar{d})] \leq \mathbb{E}_a \left[ \sup_{\mu_2: \|\mu_2 - \xi\|_2 \leq \bar{d}} \int_{a_1 \mu_2} a_2 \int_{a_1 \mu_2} \left( 1(r > v) - 1(r > a^T \mu^*) \right) d\nu(r|\bar{a}) \right] \leq \bar{\alpha} \cdot \bar{d}^2 \cdot \epsilon^2
\]
by specifying $q$ such that $\bar{\alpha} \cdot \bar{d}^2 \cdot \frac{q(1+q)}{(1-q)^2} \leq \frac{\alpha}{32}$.

(ii). If the cube $\Omega = C_{k_0}$, then by noting $\bar{\xi} = \mu^*$ and $(1 - q)^{k_0} \cdot \sqrt{\bar{m}\gamma} = \epsilon$, we have that
\[
\mathbb{E}_{(\bar{r}, \bar{a})}[g((\bar{r}, \bar{a}), \bar{\xi}, \bar{d})] \leq \mathbb{E}_{a} \left[ \sup_{\mu_2: \|\mu_2 - \bar{\xi}\|_2 \leq \bar{d}} \int_{a_1 \mu_2} a_2 \int_{a_1 \mu_2} \left( 1(r > v) - 1(r > \bar{\alpha}^T \mu^*) \right) d\nu(r|\bar{a}) \right] \leq \bar{\alpha} \cdot \bar{d}^2 \cdot \epsilon^2
\]
Therefore, for both (i) and (ii), it holds that
\[
P \left( \frac{1}{T - t} \cdot \sum_{\tau=t+1}^{T} \eta((\bar{r}_\tau, \bar{a}_\tau), (\bar{\mu}, \bar{\xi})) - \mathbb{E}_{(\bar{r}, \bar{a})}[\eta((\bar{r}, \bar{a}), \bar{\mu}, \bar{\xi})] \geq \bar{\alpha} \bar{d}^2 \epsilon^2 + \frac{\epsilon \cdot \|\bar{\mu} - \mu^*\|_2}{2} + \frac{\alpha \beta \|\bar{\mu} - \mu^*\|_2^2}{16} \right)
\]
\[
\leq P \left( \frac{1}{T - t} \cdot \sum_{\tau=t+1}^{T} g((\bar{r}_\tau, \bar{a}_\tau), (\bar{\xi}, \bar{d})) - \mathbb{E}_{(\bar{r}, \bar{a})}[g((\bar{r}, \bar{a}), (\bar{\xi}, \bar{d}))] \geq \bar{\alpha} \bar{d}^2 \epsilon^2 + \frac{\epsilon \cdot \|\bar{\mu} - \mu^*\|_2}{2} + \frac{\alpha \beta \|\bar{\mu} - \mu^*\|_2^2}{16} \right)
\]
\[
\leq P \left( \frac{1}{T - t} \cdot \sum_{\tau=t+1}^{T} g((\bar{r}_\tau, \bar{a}_\tau), (\bar{\xi}, \bar{d})) - \mathbb{E}_{(\bar{r}, \bar{a})}[g((\bar{r}, \bar{a}), (\bar{\xi}, \bar{d}))] \geq \frac{\epsilon \cdot \|\bar{\mu} - \mu^*\|_2}{2} \right)
\]
\[
\leq P \left( \frac{1}{T - t} \cdot \sum_{\tau=t+1}^{T} \eta((\bar{r}_\tau, (\bar{a}_\tau), (\bar{\xi}, \bar{d})) - \mathbb{E}_{(\bar{r}, \bar{a})}[\eta((\bar{r}, \bar{a}), (\bar{\xi}, \bar{d}))] \geq \frac{\epsilon \cdot \bar{d}}{2q} \right)
\]
We apply the Hoeffding inequality to bound the above probability for each cube $\Omega_{k_t}$ and $C_{k_0}$. It holds that
\[
g((r, a), \xi_{k_t}, d_{k_t}) \leq \bar{d} \cdot d_{k_t}, \text{ and } g((r, a), \xi_{k_0}, d_{k_0}) \leq \bar{d} \cdot d_{k_0} \forall (r, a)
\]
where $d_{k_t}$ denotes the length of the edge of the cube $\Omega_{k_t}$ and $d_{k_0}$ denotes the length of the edge of the cube $C_{k_0}$. By Hoeffding’s inequality, we have
\[
P \left( \frac{1}{T - t} \cdot \sum_{\tau=t+1}^{T} \eta((\bar{r}_\tau, (\bar{a}_\tau), (\bar{\xi}, \bar{d})) - \mathbb{E}_{(\bar{r}, \bar{a})}[\eta((\bar{r}, \bar{a}), (\bar{\xi}, \bar{d}))] \geq \frac{\epsilon \cdot \bar{d}_{k_t}}{2q} \right) \leq \exp \left( -\frac{\epsilon^2 (T - t)}{2q^2 \bar{d}^2} \right)
\]
for each cube $\Omega_{k_t}$, and
\[
P \left( \frac{1}{T - t} \cdot \sum_{\tau=t+1}^{T} \eta((\bar{r}_\tau, (\bar{a}_\tau), (\bar{\xi}, \bar{d})) - \mathbb{E}_{(\bar{r}, \bar{a})}[\eta((\bar{r}, \bar{a}), (\bar{\xi}, \bar{d}))] \geq \frac{\epsilon \cdot \bar{d}_{k_0}}{2q} \right) \leq \exp \left( -\frac{\epsilon^2 (T - t)}{2q^2 \bar{d}^2} \right)
\]
for the cube $C_{k_0}$. Note that
\[
\frac{1}{T - t} \cdot \sum_{\tau=t+1}^{T} \eta((\bar{r}_\tau, (\bar{a}_\tau), (\bar{\xi}, \bar{d})) - \mathbb{E}_{(\bar{r}, \bar{a})}[\eta((\bar{r}, \bar{a}), (\bar{\xi}, \bar{d}))] \geq \frac{\epsilon \cdot \bar{d}}{2q}
\]
implies that
\[
\frac{1}{T - t} \cdot \sum_{\tau=t+1}^{T} \eta((\bar{r}_\tau, (\bar{a}_\tau), (\bar{\xi}, \bar{d})) - \mathbb{E}_{(\bar{r}, \bar{a})}[\eta((\bar{r}, \bar{a}), (\bar{\xi}, \bar{d}))] \geq \frac{\epsilon \cdot d_{k_0}}{2q},
\]
or there exists at least one $\Omega_{k_t}$ such that
\[
\frac{1}{T - t} \cdot \sum_{\tau=t+1}^{T} \eta((\bar{r}_\tau, (\bar{a}_\tau), (\bar{\xi}, \bar{d})) - \mathbb{E}_{(\bar{r}, \bar{a})}[\eta((\bar{r}, \bar{a}), (\bar{\xi}, \bar{d}))] \geq \frac{\epsilon \cdot d_{k_t}}{2q}.
\]
Applying the union bound, our proof is completed. \(\square\)
Proof of Claim 8. Note that
\[
\eta((r, a), \hat{\xi}, \mu^*) \leq d \cdot \|\hat{\xi} - \mu^*\|_2
\]
Then, if \(\hat{\Omega} \in \{\Omega_{\theta t}\}\), we have that
\[
\eta((r, a), \hat{\xi}, \mu^*) \leq d \cdot \|\hat{\xi} - \mu^*\|_2 \leq \frac{d \cdot \|\hat{\mu} - \mu^*\|_2}{q}
\]
and if \(\hat{\Omega} = C_{\theta t}\), we have that
\[
\eta((r, a), \hat{\xi}, \mu^*) \leq d \cdot \|\hat{\xi} - \mu^*\|_2 = 0
\]
by noting \(\hat{\xi} = \xi_{\theta t} = \mu^*\). Now, for each \(\Omega_{\theta t}\), by Hoeffding’s inequality, we have
\[
P\left(\frac{1}{T-t} \sum_{r=t+1}^{T} \eta((\hat{r}, \hat{a}_r), \xi_{\theta t}, \mu^*) - \mathbb{E}_{(\hat{r}, \hat{a}_t)}[\eta((\hat{r}, \hat{a}_t), \xi_{\theta t}, \mu^*)] \geq \frac{\epsilon}{2} \cdot \|\hat{\mu} - \mu^*\|_2 \right) \leq \exp\left(-\frac{\epsilon^2 \cdot (T-t)q^2}{2d^2}\right).
\]
Our proof is completed from the union bound over all \(\Omega_{\theta t}\). \(\square\)

Proof of Theorem 2. From (5), we have
\[
\text{Myopic}_{c_t}(\pi, \hat{c}_t^*) \leq 2\hat{c} \cdot \mathbb{E}_{a_t} [\mathbb{I}(\hat{c}_t^* \geq \hat{a}_t) \cdot \text{Var}(\hat{M}_{c_t} \hat{a}_t(I_{t+1}))] + \mathbb{E}_{i_t} [G_{c_t^*}(I_t)] + 2\hat{a} \cdot \mathbb{E}_{a_t} \left[\mathbb{I}(\hat{c}_t \geq \hat{a}_t) \cdot \left(\hat{M}_{c_t} \hat{a}_t - \mathbb{E}_{i_{t+1}}[\hat{M}_{c_t} \hat{a}_t(I_{t+1})]\right)^2\right]
\]
(127)
Then, from Lemma 8 and Lemma 9, we know that
\[
\text{Myopic}_{c_t}(\pi, \hat{c}_t^*) \leq \frac{2\hat{a} \cdot \text{Var}(\hat{M}_{c_t} \hat{a}_t(I_{t+1}))}{T-t} + (2\hat{a} \cdot \text{Var}(\hat{M}_{c_t} \hat{a}_t(I_{t+1})) + 2\hat{a} \cdot \mathbb{E}_{a_t} \left[\mathbb{I}(\hat{c}_t \geq \hat{a}_t) \cdot \left(\hat{M}_{c_t} \hat{a}_t - \mathbb{E}_{i_{t+1}}[\hat{M}_{c_t} \hat{a}_t(I_{t+1})]\right)^2\right]
\]
(128)
where \(\hat{\mu}_1, \hat{\mu}_2\) are defined in (43) and \(\mu_1, \mu_2\) are defined in (47), with \(c = \hat{c}_t^*\). Moreover, from Lemma 6 and the definition of \(\hat{M}_{c_t} \hat{a}_t\) in Algorithm 3, we know
\[
\mathbb{E}_{a_t} \left[\mathbb{I}(\hat{c}_t \geq \hat{a}_t) \cdot \left(\hat{M}_{c_t} \hat{a}_t - \mathbb{E}_{i_{t+1}}[\hat{M}_{c_t} \hat{a}_t(I_{t+1})]\right)^2\right] \leq \mathbb{E}_{a_t} \left[\mathbb{I}(\hat{c}_t \geq \hat{a}_t) \cdot ((\hat{\mu}_1 - \mathbb{E}[\hat{\mu}_1])^2 + (\hat{\mu}_2 - \mathbb{E}[\hat{\mu}_2])^2)\right] \\
\leq \mathbb{E}[(\hat{\mu}_1 - \hat{\mu}_2)^2] + 2 \cdot \mathbb{E}[(\hat{\mu}_1 - \hat{\mu}_2)^2] + 2 \cdot \mathbb{E}[(\hat{\mu}_2 - \hat{\mu}_1)^2]
\]
We apply Claim 4 to bound the term \(\mathbb{E}[(\hat{\mu}_1 - \hat{\mu}_2)^2]\) and we apply the “dual convergence” established in Lemma 10 to bound \(\mathbb{E}[(\hat{\mu}_1 - \hat{\mu}_2)^2]\) and \(\mathbb{E}[\mathbb{I}(\hat{c}_t \geq \hat{a}_t) \cdot (\hat{\mu}_2 - \hat{\mu}_2)^2]\) at the order of \(O(\frac{1}{\sqrt{t}})\). Therefore, we prove our result. \(\square\)