Research Article

Evolutionary Game Theoretic Modeling and Repetition of Media Distributed Shared in P2P-Based VANET

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Received 24 April 2014; Accepted 3 June 2014; Published 23 June 2014

1. Introduction

The applications of peer-to-peer (P2P) networks have been growing at tremendous speed these past few years. Another hot area in communication networks is wireless mobile ad hoc networks, which enable distributed nodes to communicate with each other without a central access point or base station [1]. VANET is considered to be one of the most promising techniques for providing road safety and innovative mobile applications. With the proliferation of the distributed peer-to-peer (P2P) cooperative transmission technologies, P2P-based VANET has recently received a substantial amount of interest. Using the P2P services thought and at the same time, a set of methods should be designed to avoid the disadvantage of P2P system appearing in VANET. Under such a presupposition, in this paper, we study the media provisions in P2P-based VANET and present a repeated game “More Pay for More Work (RGMPMW)” incentive mechanism based on service evaluation information. We also propose evolutionary game-based veracity (EGV) game model which exploits evolutionary game to guarantee the multimedia service share veracity of all vehicles in VANET. In addition, we provide extensive simulation results that demonstrate the effectiveness of our proposed schemes.
Some incentive mechanisms had been studied in P2P system and VANET. Reference [16] studied a repeated game-based incentive mechanism of bandwidth employing a trusted third party to record peers’ contribution in each round of P2P system. Reference [17] proposed an incentive mechanism based on coalitional games, which makes the vehicle user select better route, improves the ratio of delivery, and reduces the delay. Reference [18] uses Reed-Solomon codes (RS-codes) to construct our incentive scheme and enhance its security by introducing one discrete logarithm representation problem to guarantee the vehicles cooperation in VANET. However, the incentive mechanisms in these papers studied either not apply to VANET or stimulate forwarding with completed compute. Unlike these, in this paper, under the media provision in P2P-based VANET, we exploit little feedback information to present a RGMPMW incentive mechanism to stimulate vehicles share their multimedia in V2V communication.

VANET is a kind of autonomic networks, and the vehicles in VANET are limited rationality at most cases [19]. Under the action of RGMPMW incentive mechanism, vehicles may exaggerate their contributions to get the most profits. Here, the exaggeration of service contribution contains the following: (a) vehicles exaggerate their service contributions to get most profits; (b) actions of malicious vehicles: exaggerating their service contributions intentionally to reduce the profits of competitors. Reference [16] solved the honest problem using trusted third party. However, in VANET, each vehicle is an independent node. However, in VANET, each vehicle is equipped with a node. Nodes from different vehicles can communicate, but each node manages itself and is in charge of its own service contribution in VANET. Guaranteeing the veracity of this autonomous network is a challenging problem because there is no central and trusted manager to protect the whole network. Thus, only distributed mechanisms are viable in autonomous networks such as VANET. Evolutionary game theory is used to the model where players have limited rationality in the game; thus it is suitable to analysis vehicles with limited rationality, so we can use evolutionary game theory to investigate model. The application of evolutionary game in network is popular in recent years [20–25]. Therefore, in this paper, we propose EGV game model to prevent vehicles are not veracity, improve the stability of system, and ensure the validity of the mechanism.

Therefore, considering the specific characteristic of multimedia in P2P-based VANET, this paper mainly studies the following several aspects.

1. In P2P-based VANET, we present a RGMPMW incentive mechanism based on the information of services evaluation. Repeated game is exploited to accurately evaluate the service contribution of each vehicle in every game stage, stimulating vehicles to share their multimedia in V2V communication.

2. Based on evolutionary game, propose EGV game model to expand RGMPMW incentive mechanism of media. Regarding the several vehicles that request the same supplier in the VANET as a population, prevent the mendacious service share of vehicles.
efficiently and guarantee service contribution veracity of vehicles in each stage game by evolving of single population.

The target of this paper is to present a RGMPMW incentive mechanism based on the information of services evaluation to stimulate vehicles sharing their multimedia based on evolutionary game, propose EGV game model to prevent the mendacious service share of vehicles efficiently, and guarantee service contribution veracity of vehicles in each stage game by evolving of single population in P2P-based VANET.

The rest of the paper is organized as follows. Related work on evolutionary game and repeated game are presented in Section 2. Section 3 is system model. Section 4 describes feedback-based RGMPMW incentives scheme. Evolutionary game-based model for veracity of vehicles is reported in Section 5. Section 6 studies the simulation and analysis. We conclude our paper and propose perspectives in Section 7.

2. Related Work

Game theory is a mathematical theory and method of researching the phenomenon with wrestled or competitive property, and, under certain restricted conditions, the participants can implement strategies to other players, to make sure that the profits of participants reach a final balanced state.

2.1. Evolutionary Game. Classical game theory generally assumes that each individual acts so as to maximize its utility, assuming that the others do the same and understanding completely the possible utility payoffs of their interactions. An alternative perspective from evolutionary game theory assumes that individuals will copy the behavior of others who obtain a higher utility [22]. Evolutionary game has two core theories: evolutionary stable strategy and replicator dynamics equation. Evolutionary stable strategy emphasizes the process of how a dynamic evolution system reaches a steady state, and evolutionary stable strategy \( x^* \) needs to meet the following two conditions: first, \( x^* = \Omega(x^*) = 0 \); second, \( \Omega^I(x^*) < 0^6 \). Replicator dynamic equation describes a dynamic differential equation on time parameter \( t \), taking the number of frequency variation. Replicator dynamic equation is expressed by \( x'_i = (\phi(x_i, x) - \phi(x, x))x_i \), where \( x_i(t) \) is the proportion of participants who take pure strategy, \( \phi(x_i, x) \) is the fitness of strategy \( i \), and \( \phi(x, x) \) is the average fitness of strategy \( i \) [26].

The following describes the standard set of evolutionary game theory [27].

(1) There is a group of users, and the number of the users in the group is very large.

(2) Assuming that pure strategy or action exists. Each member in the group chooses the strategy from the same strategy set \( A = \{1, 2, \ldots, I\} \).

(3) Let \( M := \{y_1, \ldots, y_I\} \ | \ y_j \geq 0, \sum_{j=1}^I y_j = 1\} \) be the possibility distribution set on pure strategy set \( I \). \( M \) can be interpreted as a mixed strategy set. In fact, assuming that picking the user randomly from the group, the possibility that the user with marked sign meets the user with strategy \( y \) is \( y_j \). After a few game processes, participators who use strategy \( j \) is equivalent to the face using a mixed strategy of participators.

Evolutionary game in network application is very popular and a currently [20–25, 28] used evolutionary game to solve the problem of uncooperative two-hop routing message forwarding control in DTN (delay tolerant network). Reference [21] used the mathematical framework of game theory and evolutionary game theory for modeling, clearly demonstrating the connection of cooperation, trust, privacy, and security in a multihop network, in order to prevent the fake nodes, prevent misuse, detect anomalies, and to protect users’ privacy. Reference [22] studied the resource allocation and the establishment of distribution tree. Assuming that the allocation of video is based on a multiple description encoder, then using evolutionary game can significantly affect the scalability, fast adaptation twist, high degree of node cooperation, and the autonomous of distributed node tree network. Reference [23] studied an energy management of stochastic evolutionary game in distributed aloha network. Each participator may be in different states and can be involved in the same local interaction during its lifetime. Its action decides not only effectiveness, but also the transmission probability and the time of his life. Reference [24] analyzed the problem of network dynamic selection with evolutionary game in heterogeneous wireless networks. It can guarantee the network performance in the case of multipopulation evolutionary game competition. References [25] studied the evolutionary game based on the credibility to forward data safely and deny the service attacks, making sure that the maximum numbers of nodes in the network forward cooperatively in the autonomous ad hoc and sensor network.

In P2P-based VANET, vehicles move at high speed, and instant contribution is more important; therefore in every request slot \( t \), each vehicle needs to decide how many media services to share with the request vehicle. This decision will affect its profit in the next stage. Thus, it is reasonable to model the action of vehicles as repeated game.

2.2. Repeated Game. Repeated game means the same structure game is repeated for many times, where each game is called “stage game.” Repeated game is an important part of a dynamic game, and it can repeat complete information or incomplete information. When the game is executed only once, each participant just cares about one-time payment. But if the game is repeated, participants may tend to get long-term profit instead of immediate profit; therefore they will choose a different equilibrium strategy. Thus, the repeat number will affect the outcome of equilibrium game. Repeated game has three basic characteristics. (a) In the repeated game stage, there is no “substance” connection between games that is to say that the previous stage of the game will not change the construction of the next stage. (b) In every stage of the repeated game, all participators can observe
the history of the game. (c) Total profit of participators is the discounted value sum or weighted average number of all stages’ profit.

Repeated game in network application has also been studied widely [16, 29–32]. Reference [16] proposed server differentiation incentives for P2P streaming system based on the immediate profit of nodes. At the same time, it designed a repeated game model to analyze how much should every node contribute in every round in this incentive. Reference [29] used repeated game theory in FiWi access network and applied effective quality service routing mechanisms and scheduling policy into practical application. Balance strategy guaranteed the quality of service of FiWi wireless network. Reference [30] used repeated game model to establish a limited punishment mechanism to enforce selfish nodes to be unselfish, preventing cheating to save energy. Reference [31] checked some basic important properties of a routing protocol design. The importance of these attributes is autonomous participators from the underlying economic factors’ management behavior. The connected price information in an associated swap is regarded as a repeated game among the relevant participators. Reference [32] studied the multicast overlay network applications in the framework of repeated games and described a repeated game model of user behaviors, to capture the effect of short-term profit to long-term profit.

VANET is an autonomous network, and participators are groups or individuals with limited rationality. Traditional game theory methods assume that participators are entirely rational. Evolutionary game theory is a game addressing node bounded rationality specifically. Therefore, it is reasonable that we use evolutionary game as a research method in this paper.

3. System Model

In VANET, there are two main ways of communication: vehicle to vehicle (V2V) and vehicle to RSU (V2R). In the analysis of part I, we only consider V2V communication under urban scenes. We regard VANET as a network of vehicles’ set, and each vehicle is equipped with communication equipment, allowing communication based on 802.11p protocol among different vehicles. As is shown in Figure 1, nodes begin to download the initial service from RSU then when the vehicle is moving out of the range of RSU services. In order to obtain a satisfactory quality of service, V2V communication is needed in RSU communication blackout. Considering a typical P2P-based VANET multimedia services under urban scenes; at a certain time, the role of each vehicle is “Requester” (ask for service) or “Supplier” (provide service). As a requester, it may face several suppliers with same services; as a supplier, it may also face more than one requester.

The “More Pay for More Work (RGMPMW)” incentive mechanism and EGV game model are used into the urban VANET with double lane in this paper, and the incentive mechanism performs well. The models in the paper are introduced as follows.

3.1. The Model of Vehicles Encounter. VANET is a network collection of vehicles; \( V = \{1, 2, \ldots, i\} \). Each vehicle can join or leave the network at any time. The running speed of vehicle is \( v_i, i = \{1, 2, \ldots, |V|\} \) m/s. We use \( d_{i,j} \in (1, 5000) \) to express the distance between vehicle \( i \) and vehicle \( j \), and the chance of encounter between vehicles is independent with no interfering. Here we define vehicles that meet within each transmission range. Assuming the transmission range of the every vehicle is 250 m. The probability that vehicle \( i \) and vehicle \( j \) meet is expressed as \( q_{i,j} \):

\[
q_{i,j} = \begin{cases} 
\frac{v_i - v_j}{0.05 \cdot d_{i,j}}, & \text{same running direction, vehicle } i \text{ is behind vehicle } j, \ v_i > v_j; \\
0, & \text{leaving in opposite direction or same running direction, vehicle } i \text{ is behind vehicle } j, \ v_i < v_j; \\
1, & \text{move in opposite direction, } d_{i,j} \in (250, 5000) \text{ or } d_{i,j} \in (1, 250). 
\end{cases}
\]

(1)

As is shown in Figure 2(a), in the period \( t - 1 \), vehicle 2 has the media service which vehicle 1, vehicle 5, and vehicle 6 need; then they make a service request to vehicle 2. At this time, node 2 can be regarded as a similar manager of nodes 1, 2, 5, and 6. In the next period \( t \), there are two vehicle nodes 7 and 8 added into network. At this time vehicles 5 and 6 run out of the communication range of vehicle 2; therefore they make media request from vehicle 7. Vehicle 1 and vehicle 2 are running in the same direction, so they can still continue to keep connection. At this time, vehicle 2 loses manage capabilities of vehicles 5 and 6, and vehicle 7 becomes the current similar manager of vehicles 5 and 6; vehicle 2 becomes the similar manager of vehicles 1, 2, 4, and 8. With the movement of vehicles, every vehicle’s similar manager is changing, in other words, the connected time between vehicles is not fixed. Figure 2 is just a simplified model, only showing a partial service communication. In fact, every vehicle is a similar manager of a group. In this paper, we only consider the situation of one provider and several requesters.

Because of the mobility of vehicles in VANET, it is impossible to choose a fixed reliable third-party to manage a certain number of nodes. For simplicity, first we assume that (1) nodes in VANET have their own media services; (2) initially each vehicle gets a certain amount of media services from the RSU and gets initialized with a certain contribution value when entering the network; (3) the media service among vehicles is instantaneous, and there is no time limit; (4) at the time every vehicle enter the network, it is given
a unique real identity; (5) the transmission range of vehicles is the same.

3.2. Media Service Model. In this paper, we divide multimedia services into four kinds based on the type and popularity of the media: (1) critical urgent media services, such as road hazard information and highway information, defined as $P_1(t) = 0.9$; (2) delay-sensitive services [33], such as video conference and video service, defined as $P_2(t) = 0.7$; (3) constantly completed multimedia services, like music entertainment, defined as $P_3(t) = 0.5$; (4) life services, such as restaurant, hotel information service, defined $P_4(t) = 0.2$, where $P_k(t)$ represents the popularity of the media provided by providers in the current stage of the game, $P_k(t) \in [0,1]$.

When requesters ask for media service from provider, they will give their shared value of the previous phase to the provider. Since vehicles have a strong ability of computing, provider will determine the allocation of resources based on the shared value of all requesters. Meanwhile, at the end of the request, requesters will broadcast information to all nodes in the network, such as information about media service, provider, and value (Section 4 will describe evaluation of media service in detail). Therefore, all nodes in network will store a list: record ID and shared value of all nodes in the network. For nodes that just enter the network, they will get a unique ID, and then when they meet first node, they will copy all records from it to complete information. All nodes update records in each time slot.

4. Offer-Based RGPMMW Incentives Scheme

Since vehicle nodes in VANET are naturally selfish, they will not go to contribute their resources to other peer nodes without motive. Therefore, we need to design an incentive mechanism to encourage the contribution of nodes [17].

In the P2P-based VANET, design of incentive mechanism should consider instant contribution of vehicles. Taking real-time requirement of media streams into account, the vehicles are more strictly required to share their resources in every round; otherwise, the requester may not receive the data before the playback time. When implementing incentive mechanism to nodes, contribution of current time period is more important than historical contribution. Repeated game keeping encouraging nodes to contribute includes a lot of repeated game stages of participators. In each stage, the decisions of participators all depend on their payment. An action can be determined by one participator, giving it the highest payment. Therefore, in the repeated game, when participators decide what strategy to take, they must care about current and future payment [29]. This paper proposes a RGPMMW incentive mechanism based on similar managers.

In this incentive mechanism, we define a noun “shared contribution value,” representing the contribution of a node made in a game stage. It is related to bandwidth, popularity of the media, and amount of providers of nodes’ contribution. In the sharing mechanism, “shared contribution value” of nodes is evaluated by upload/download behavior in a previous stage. Each node broadcasts evaluation of service’s popularity and importance before the end of the game. In the same stage, every provider is a similar manager, and it decides requesters’ profit in current stage based on “shared contribution value” of previous stage provided by requesters.

Therefore, the “shared contribution value” of a node kin stage $t$ is the sum of feedback information provided by all nodes that received the media service of $k$. That is,

$$V_k(t) = \sum_{i=1}^{n} C_i(t) \ast P_i(t).$$

For simplicity, we assume that the system model in this paper is that every vehicle node in a slot $t$ provides only one kind of media service. Thus, the “shared contribution value” of a node $k$ in slot $t$ can be simplified as

$$V_k(t) = C_k(t) \ast P_k(t),$$

where $V_k(t)$ represents “shared contribution value,” $C_k(t)$ represents bandwidth contributed by $S$, and $P_k(t)$ represents

**Figure 2:** The model of vehicles encounter.
Figure 2 shows the system model of communications among vehicles under urban scenes in VANET. A vehicle may be both a service provider and a service requester in a stage of the game. However, in the incentive mechanism proposed in this paper, we are most concerned about the "shared contribution value." Here we only consider a simple scenario: one provider corresponds to several requesters, as shown in Figure 3.

In vehicle request model, when a vehicle asks for media services, it will give its "shared contribution value" in previous stage to provider, and provider will determine the allocation of resources based on the shared value of all requesters. We define the resource assigned from provider $S$ to a requester in time $t$ which is

$$G_{s,i}(t) = q_{i,s} \cdot G_s \cdot \frac{V_i(t-1)}{V_i(t-1) + V_{-i}(t-1)},$$

(4)

where $q_{i,s}$ is the meet probability of two vehicles, and it is related to their the running speed and distance; $G_s$ is the contribution set by provider $S$ in current stage; $V_i(t-1)$ is the "shared contribution value" of requester $i$ in previous stage; $V_{-i}(t-1)$ is the sum "shared contribution value" of all requesters except requester $i$. Therefore, the profit function of each node is the difference between the service obtained by provider and the total service it provides for other requesters:

$$u_i(t) = G_{s,i}(t) - V_i(t).$$

(5)

As vehicle node is autonomous, the service media type is decided by each node. In order to obtain greater payoff, each node tends to choose high popularity of media services. We set $G_s = V_i(t-1)$. Then, the utility function of requestor $i$ can be rewritten as

$$u_i(t) = q_{i,s} \cdot V_i(t-1) \cdot \frac{V_i(t-1)}{V_i(t-1) + V_{-i}(t-1)} - V_i(t).$$

(6)

In addition, we define the total utility of node $i$ in the service time as follows:

$$u(i) = \sum_{t=1}^{\infty} \delta (1 - p)^{t-1} u_i(t).$$

(7)

Here $\delta \in (0, 1]$ is the discount factor, and it can be regarded as a node's patience for the subsequent game. The greater the value is, the more patient node is. On the contrary, the node will pay more attention to the current earnings. In the infinite repeated game, every participant does not know when the game will end. So, we assume that the probability of the end of the game is $p$.

The implementation steps of RGMPMW incentive mechanism based on similar management are shown in Algorithm 1.

We can get from formulas (6) and (7) the following:

$$u(i) = \sum_{t=1}^{\infty} \delta (1 - p)^{t-1}$$

$$* \left\{ q_{i,s} \cdot V_i(t-1) \cdot \frac{V_i(t-1)}{V_i(t-1) + V_{-i}(t-1)} - V_i(t) \right\}.$$ 

(8)

The goal of the node $i$ is to maximize $u(i)$. First, we can get the formula (9) by a series of deformation for the formula (7):

$$u(i) = \frac{V_i(0)}{V_i(0) + V_{-i}(0)} + \sum_{t=1}^{\infty} \delta (1 - p)^{t-1}$$

$$* \left\{ (\delta (1 - p)) \cdot q_{i,s} \cdot V_i(t) \right\}$$

(9)

$$* \left\{ \frac{V_i(t)}{V_i(t) + V_{-i}(t)} - V_i(t) \right\},$$

where $V_i(0)$ is the initialization "shared contribution value," when the node $k$ comes into the network at the beginning. So, $V_i(0)/(V_i(0) + V_{-i}(0))$ is a constant. After deformation, each item is independent in the sum. Therefore, we can make the sum maximized by maximizing each item.

We set

$$d([\delta (1 - p)] \cdot q_{i,s} \cdot V_i(t) \cdot (V_i(t) / (V_i(t) + V_{-i}(t)))) - V_i(t))$$

$$dV_i(t) = 0.$$ 

(10)

That is,

$$[\delta (1 - p)] \cdot q_{i,s} \cdot V_i(t) \cdot \frac{V_{-i}(t)}{[V_i(t) + V_{-i}(t)]^2} - 1 = 0.$$ 

(11)
Input: initialize the $V_i(t)$  
Output: the optimum $V_i(t)$ for the max profit $u(i)$
Procedure RGMPMW
for $t = 0$ to $\infty$
  if (vehicle $i$ begins to request media service)
    if (vehicle $i$ meets vehicle $S$ & vehicle $S$ has the media service)
      Vehicle $i$ provide the $V_i(t-1)$ for the vehicle $S$
      and requests media;
      Vehicle $S$ computes how much to share for vehicle $i$:
      $G_{s,i}(t) = q_{i,s} \cdot G_s \cdot \frac{V_i(t-1)}{V_i(t-1) + V_{-i}(t-1)}$
      Vehicle $i$ games with the vehicle $-i$;
      if (exist a vehicle $j: u(j) > u(i)$)
        Vehicle $i$ change its strategy;
    endif
  endif
endfor
end

Algorithm 1: RGMPMW incentive mechanism.

Then the optimal solution $V_i^*(t)$ of $V_i(t)$ is

$$V_i^*(t) = \sqrt{\left[ \delta \left(1 - p \right) \right] \cdot q_{i,s} \cdot V_s(t) \cdot V_{-i}(t) - V_{-i}^*(t)}.$$  (12)

We know that in the condition of NE (Nash equilibrium), the following formula is true for each node:

$$V_i^*(t) = \sqrt{\left[ \delta \left(1 - p \right) \right] \cdot q_{i,s} \cdot V_s(t) \cdot V_{-i}(t) - V_{-i}^*(t)}.$$  (13)

Therefore, we have

$$V_{-i}^*(t) = (n - 1) \cdot V_i^*(t),$$  (14)

where $n$ is the number of vehicles requesting node $S$.

Putting the formula (14) to the formula (12), we can get

$$V_i^*(t) = \frac{(n - 1) \cdot \left[ \delta \left(1 - p \right) \right] \cdot q_{i,s} \cdot V_s(t)}{n^2}$$

$$= \frac{(n - 1) \cdot \left[ \delta \left(1 - p \right) \right]}{n^2} \cdot C_s(t) \cdot P_s(t) \cdot q_{i,s}.\quad (15)$$

The payoff of the node will be maximized, when the equation above equation is established.

5. Evolutionary Game Model for Veracity of Vehicles

In Section 4, the function of RGMPMW incentives scheme based on information is when the node is rational; it will actively share its media resource to gain more payoffs. But in VANET which is autonomic network, it is not practical for the node to be completely rational. In the process of each game, if each requestor’s “shared contribution value” is stable at present, namely, node “shared contribution value” fluctuation is small, he will get unexpected payoff when one requester exaggerates his previous contribution; or there are attacks of malicious nodes, which exaggerate their contribution deliberately and make the payoffs low. So we should study node bounded rationality in VANET and the situation that the nodes do not trust each other. In this section, we present an EGV game model by using a game theory which can be applied to the node bounded rationality, which can prevent the node exaggerating its “shared contribution value” to gain extra payoff or malicious attacks and to guarantee the authenticity of all the nodes.

5.1. Structure and Solution of Evolutionary Game. In this paper, we set the vehicles which request for the same vehicle as an evolutionary group. Researching on evolutionary game theory, a mutation of disadvantage group is the vehicles which exaggerate their own shared services for more payoffs and reduce the payoff of the other competitors. After a long evolution, disadvantage group will be eliminated and vehicles’ real “shared contribution value” will be guaranteed. We know that in the real network the exaggerated nodes will benefit more, because it means that in the case the other nodes are real, the exaggerated nodes will get more in the next round of the game.

In VANET based on P2P, vehicles provide service for each other. In each stage in the game, all the vehicles will broadcast their gains. So all the vehicles in the network will receive the broadcast information and accumulate the value according to the identity of the vehicles. At the end of the stage game, each vehicle records the stage game information (vehicle identification, total service) of all the vehicles in the network. Therefore, the vehicles will refuse to provide service when the nodes choose “exaggerator” according to their record.

In evolutionary game, the game of the participants is two random vehicle nodes. Suppose in the whole population
that the population strategy set is \{real, exaggerated\}. If the two participants \(I\) and \(j\) are both real, their payoffs are \(u(i)\) and \(u(j)\); if participants exaggerate their services, it will be punished and the provider refuses to provide them services; the real party will get the rewards \(\alpha f\), where \(f\) is the reward unit and \(\alpha\) is the strength of the reward. Therefore, the pay-off matrix is as in Table I.

We define that \(y_0(t)\) is the number of nodes choosing the “real” strategy, and \(y_1(t)\) is the number of nodes choosing the “exaggerator” strategy. Their relation is

\[
y(t) = y_0(t) + y_1(t).
\]

We set \(x(t) = y_0(t)/y(t)\) on behalf of the proportion of the peer following the strategy “real”; then proportion of the peer following the strategy “exaggerate” is \(1 - x(t)\).

According to the game matrix, the payoff of game parties choosing the real strategy is

\[
U_k^V = u(k) \ast x(t) + [u(k) + \alpha f] \ast [1 - x(t)]
\]

\[
= u(k) + [1 - x(t)] \alpha f.
\]

The payoff of choosing exaggerated strategy is

\[
U_k^A = -V_k(t).
\]

The average payoff is

\[
U_k^A = x(t) \ast U_k^V + [1 - x(t)] U_k^N
\]

\[
= \{u(k) + [1 - x(t)] \alpha f\} x(t) - V_k(t) [1 - x(t)].
\]

The replication dynamic below indicates how evolution makes dynamic change; in particular, it can be converted to the equilibrium dynamically by replication dynamic. Replicator dynamic describes a population evolution process with multiple strategies. Each individual in the population obeys the following imitation rules: after studying the individual choose the strategy getting more benefit.

We assume that each stage game begins from \(kt, k \in N\), and ends at \((k + 1)t, k \in N\). The average payoff of the node is related to game rivals. Suppose in a very small time interval \(\epsilon\) that only the \(\epsilon\) part participates in the game. So in time \(t + \epsilon\), the nodes’ average payoff for adopting strategy \(i\) can be expressed as [20]

\[
y_i(t + \epsilon) = (1 - \epsilon) y_i(t) + \epsilon y_j(t) U_i(t); \quad i = 0, 1,
\]

where \(U_0(t) = U_k^V\) and \(U_1(t) = U_k^N\). Therefore, in the whole network, we have

\[
y(t + \epsilon) = (1 - \epsilon) y(t) + \epsilon y(t) \bar{U}(t),
\]

where \(\bar{U}(t) = U_k^A\). Divided (21) by (20). We can get a frequency equation for the strategy of “real”:

\[
x(t + \epsilon) - x(t) = \frac{x(t) \left[U_0(t) - \bar{U}(t)\right]}{1 - \epsilon + \epsilon \bar{U}(t)}.
\]

Then, we divide \(\epsilon\) at both sides of the equation and get

\[
x(t + \epsilon) - x(t) = \frac{x(t) \left[U_0(t) - \bar{U}(t)\right]}{1 - \epsilon + \epsilon \bar{U}(t)}.
\]

When \(\lim \epsilon \to 0\), we have

\[
\frac{dx(t)}{dt} = x(t) \left[U_0(t) - \bar{U}(t)\right].
\]

That is the Dynamic replication equation of game participant \(k\) is

\[
\frac{dx(t)}{dt} = x(t) \left(U_k^R - U_k^A\right)
\]

\[
= x(t) \{u(k) + [1 - x(t)] \alpha f
\]

\[
- [u(k) + [1 - x(t)] \alpha f] x(t)
\]

\[
+ V_k(t) [1 - x(t)]\}

\[
= \{u(k) + [1 - x(t)] \alpha f + V_k(t)\} [x(t) - x^2(t)].
\]

We set \(F(x) = dx/dt\), so

\[
F(x) = \frac{x(t + 1) - x(t)}{\Delta t}
\]

\[
= \{u(k) + [1 - x(t)] \alpha f + V_k(t)\} \left(x(t) - x^2(t)\right).
\]

According to the first condition ESS (evolutionary stable strategy) meeting, we make \(dx/dt = 0\); that is,

\[
\{u(k) + [1 - x(t)] \alpha f + V_k(t)\} \left(x(t) - x^2(t)\right) = 0.
\]

The solution is \(x_1(t) = (u(k) + V_k(t))/\alpha f + 1, x_2(t) = 0, x_3(t) = 1\).

5.2 Stability Analysis. The above three conditions of solutions are not all ESS. We need according to the second condition ESS meeting to analyze the stability.

**Theorem 1.** In EGV game model, there is only an evolutionary stable strategy of ESS.

**Proof.** According to the second condition ESS meeting, we know that in the ESS \(F(x)\) meet the conditions are

\[
F(x^*) = 0,
\]

\[
F'(x^*) < 0.
\]
Therefore, we have the analysis as follows.

(1) According to the introduction of RGMPMW incentive mechanism in Section 4, \( u(k) > 0, V_k(t) > 0 \), because it is the reward of real participants, \((u(k) + V_k(t))/\alpha f > 0\). And because \( x \) is the ratio of choosing real; that is, \( x(t) \in [0,1] \), \( x(t) \) cannot equal to \((u(k) + V_k(t))/\alpha f + 1\).

(2) Next we analyze the case when \( x_2 = 0, x_3 = 1 \). According to the analysis of (1) we can get \( u(k) + (1 - x)\alpha f + V_k(t) > 0 \). Therefore, replication dynamic evolution graph is as in Figure 11.

Assuming that there are \( \eta \) proportion of players in the game deviating from the strategy “real” and select the “exaggerated,” there are

\[
U_k^V = (1 - \eta) * u(k) + \eta * [u(k) + \alpha f] = u(k) + \eta * \alpha f,
\]

\[
U_k^{NV} = -V_k(t),
\]

\[
U_k^A = (1 - \eta) * U_k^V + \eta * U_k^{NV}
\]

\[
= [u(k) + \eta * \alpha f] (1 - \eta) - \eta * V_k(t),
\]

\[
U_k^V = u(k) + \eta * \alpha f > 0 > U_k^{NV}.
\] (29)

Therefore, \( x(t)_3 = 1 \) is the evolution stable strategy ESS.

Assuming that there are \( \eta \) proportion of players in the game deviating from the strategy “exaggerated” and select the “real,” there are

\[
U_k^V = \eta * u(k) + (1 - \eta) * [u(k) + \alpha f]
\]

\[
= u(k) + (1 - \eta) * \alpha f,
\]

\[
U_k^{NV} = -V_k(t),
\]

\[
U_k^A = \eta * U_k^V + (1 - \eta) * U_k^{NV}
\]

\[
= [u(k) + (1 - \eta) * \alpha f] \eta (1 - \eta) - (1 - \eta) * V_k(t),
\]

\[
U_k^V = u(k) + (1 - \eta) * \alpha f > U_k^{NV}.
\] (30)

So \( x(t)_2 = 0 \) is not the evolutionary stable strategy.

In conclusion, in the EGV game model, the ESS is only \( x^* (t) = 1 \).

The proving is over. \( \square \)

The above analysis of stability shows that whether the population of participants choose real or exaggerated, after a period of evolution, all the participants will choose the pure strategy—real. The proposed game model EGV ensures the authenticity of all participants.

5.3. Influence Factor Analysis of ESS. According to the analysis in Section 4, the benefits of a node \( k \) are as follows:

\[
u(k) = \frac{V_i(0)}{V_k(0) + V_{-k}(0)} + \sum_{t=1}^{\infty} [\delta (1 - p)]^{t-1} \]

\[\left\{ [\delta (1 - p)] q_{i,s} V_i(t) \right\} \]

\[
\frac{V_k(t)}{V_k(t) + V_{-k}(t)} - V_k(t). \] (31)

We set \( V_i(0)/(V_k(0) + V_{-k}(0)) = u(0) \); then we get the optimal solution:

\[
V_i^*(t) = \frac{(n - 1) \cdot [\delta (1 - p)]}{n^2} \cdot C_s(t) \cdot P_s(t) \cdot q_{i,s}. \] (32)

Setting it into formula (31), we can get

\[
u(k) = u(0) + \sum_{t=1}^{\infty} [\delta (1 - p)]^t \cdot \frac{C_s^2(t) \cdot P_s^2(t) \cdot q_{i,s}^2}{n^2}. \] (33)

When \( n \) is large enough, the profit is \( u(0) \). This is because, there are many vehicles competing for resources; their revenue is negligible and the additional income is essentially zero.

Reformatting the formula (25) and putting it into the \( u(k) \) provide the following:

\[
x(t + 1) = x(t) + 0.1x(t) \cdot [1 - x(t)] \]

\[\left\{ u(k) + [1 - x(t)] \alpha f + V_k(t) \right\} \]

\[
x(t) + 0.1x(t) \cdot [1 - x(t)] \]

\[\left\{ u(0) + \sum_{t=1}^{\infty} [\delta (1 - p)]^t \cdot \frac{C_s^2(t) \cdot P_s^2(t) \cdot q_{i,s}^2}{n^2} \right\}. \] (34)

Therefore, the impaction factors on ESS that we can get from formula (34) are as follows:

(1) the reward of choose real: \( \alpha \); 
(2) the number of participants: \( n \); 
(3) the multimedia types, that is, the “shared contribution value” of node \( k \) at the current stage: \( V_k(t) = C_k(t) \cdot P_k(t) \); 
(4) the encounter probability of the vehicles: \( q_{i,s} \); 
(5) the concrete analysis is in simulation part.
6. Simulation and Analysis

6.1. Simulation Settings. The system parameters of simulation settings are shown in Table 2. The vehicle is random distribution. Vehicles that provide service probability in a slot t are living service: to complete the media: delay sensitive services: the emergency information service = 1:1:1:2.

| Parameter                  | Value      |
|----------------------------|------------|
| The coverage of vehicle    | 250 m      |
| The speed of vehicle       | $v_i \in (5, 16)$ m/s |
| The distance between vehicles | $d_{ij} \in (1, 5000)$ m |
| The discount factor        | $\delta = 0.98$ |
| The game ended probability | $P = 0.2$ |

(2) Correct and Effective Incentive Mechanism. Figure 5 shows the effectiveness of the RGPM MW incentive mechanism: after a period of incentive, the node utility will reach a maximum. Node will increase their "shared contribution value" for its benefit. We design the RGMWMP incentive mechanism to make the nodes share their resources as much as possible positively, that is, to make the node "shared contribution value" increase. It can be seen from the above two figures that there is a game equilibrium state which makes the benefit reach the maximum. The corresponding "shared contribution value" of bigger one of two $U(k)$ from Figures 5(a) and 5(b) is the same as the stable one from Figure 4(b) when $n = 3, n = 5$, respectively. It indicates the correctness and effectiveness of the incentive mechanism of RGPM MW that we design.
6.3. EGV Game Model

(1) Validity Analysis. Figure 6 shows that, when the vehicle group has 50% vehicles select exaggeration, after a period of evolution, they will be eliminated. All the vehicles will select “real”. The results show that, in the vehicle in the group use the EGV game model can obtain satisfactory results. It proves that the EGV game model we proposed is effective.

(2) Analysis of Influence Factors

(a) Initial Value $x(0)$. As shown in Figure 7, in the vehicle group, the larger “real” ratio of vehicles is at the beginning stage of EGV game, the faster group ESS reaches. Because, if more vehicles select “real” in groups, then when the vehicles selecting “exaggerator” select game opponent, the probability of selecting real vehicle is relatively large. In the game learning process, the exaggerative will become “real.” Therefore, the vehicles group will quickly change their strategies and reach the ESS faster.

(b) Incentive Strength $\alpha$. Consider $\alpha = 1$ (hotel restaurant service) $\alpha = 5$ (immediate service); $\alpha = 8$ (delay sensitive services) $\alpha = 12$ (emergency media service).

Figure 8 shows, when the incentive strength is greater, the group tends to the ESS quicker. The reason is that the incentive strength is greater and can lead the vehicle to have higher incentives. In the dynamic evolution process, there
will be more participants who choose strategies to maximize their own real earnings.

(c) Effects of \( N \) Number of Participants. When the number of vehicles in group becomes bigger, that is to say the more number of vehicles to exaggerate, then, in the EGV game, it will converge more slowly to ESS, as shown in Figure 9. But when the number of vehicles involved in the game reaches a certain amount in the group, there was no change in convergence speed. Because of the increasing number of participants, the learning process become very widely. When the number of participants increased to a certain extent, the evolution convergence speed is no longer affected by the number of participants.

(d) Multimedia Types. Set bandwidth: \( C = 5 \). We put the multimedia service divided into four types: (1) the key emergency media services, such as “Danger! Information” and highway information, \( P_i(t) = 0.9 \); (2) delay sensitive services, such as video conference and video service, \( P_i(t) = 0.7 \); (3) immediate complete multimedia services, such as music and entertainment, \( P_i(t) = 0.5 \); (4) the life service, such as restaurants, hotel information, \( P_i(t) = 0.2 \).

As shown in Figure 10, the sharing of multimedia services is more popular; the vehicles tend to stability more quickly. Because the multimedia types not only affect the real vehicle incentives, but also affect the vehicle "shared contribution value," multimedia is more popular, and vehicles "share contribution value" is bigger, which can also give the option of the real vehicle reward greater efforts. Thus, the vehicle shares more multimedia popular, can incentive mechanism under the RGMPMW faster to achieve stability, and the vehicles will get more reward. Group will arrive at ESS steady state, as shown in Figure 10. That the vehicles will share the popular media more actively, making emergency news media service timely diffusion in VANET, which is the result we want.
7. Conclusions and Perspectives

In this paper, we studied media services in P2P-based VANET, where all vehicles are regarded as individuals with limited rationality. We proposed “More Pay for More Work (RGMPMW)” incentive mechanism to encourage vehicle nodes to share resources and studied evolutionary game to guarantee the service share veracity of all vehicles. With “shared contribution value,” RGMPMW incentive mechanism accurately evaluated the contribution of each node based on similar manager. Then as expansion to RGMPMW incentive mechanism, EGV game model had been studied to prevent the mendacious service share of vehicles efficiently. The simulation results proved RGMPMW incentive mechanism and EGV game model are correct and effective in VANET. In particular, the analysis of factors ESS shows that the fewer the number of participants is, the more urgent multimedia services are, and the faster the ESS will reach. At the same time, the proposed mechanism can be well adapted to the V2V communication with high mobility and fast topology changes.

We only considered the most simple P2P-based VANET scene, that is, one provider to several requesters. In future work, we will study evolutionary game in more complicated scene of several-to-several, including variations between nodes and unequal connection probabilities in multiple groups.

Conflict of Interests

The authors declare that there is no conflict of interests regarding the publication of this paper.

Acknowledgments

This work is supported by National Natural Science Foundation of China (61376201), the Scientific Research Foundation for the Returned Overseas Chinese Scholars (45), Liaoning Provincial Natural Science Foundation of China (2013020019), and High-Tech 863 Program (no. 2012AA111902).

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