AN ALGORITHM TO COMPUTE ROTATION NUMBERS IN THE CIRCLE

LLUÍS ALSEDA
Departament de Matemàtiques, Edifici Cc, Universitat Autònoma de Barcelona, 08193 Bellaterra (Barcelona), Spain
Centre de Recerca Matemàtica, Campus de Bellaterra, Edifici Cc, Universitat Autònoma de Barcelona, 08193 Bellaterra (Barcelona), Spain

SALVADOR BORRÓS-CULLELL
Departament de Matemàtiques, Edifici Cc, Universitat Autònoma de Barcelona, 08193 Bellaterra (Barcelona), Spain

Abstract. In this article we present an efficient algorithm to compute rotation intervals of circle maps of degree one. It is based on the computation of the rotation number of a monotone circle map of degree one with a constant section. The main strength of this algorithm is that it computes exactly the rotation interval of a natural subclass of the continuous non-invertible degree one circle maps.

We also compare our algorithm with other existing ones by plotting the Devil’s Staircase of a one-parameter family of maps and the Arnold Tongues and rotation intervals of some special non-differentiable families, most of which were out of the reach of the existing algorithms that were centred around differentiable maps.

1. Introduction

The rotation interval plays an important role in combinatorial dynamics. For example Misiurewicz’s Theorem [9] links the set of periods of a continuous lifting $F$ of degree one to the set $M := \{ n \in \mathbb{N} : \frac{k}{n} \in \text{Rot}(F) \text{ for some integer } k \}$, where $\text{Rot}(F)$ denotes the rotation interval of $F$. Moreover, it is natural to compute lower bounds of the topological entropy depending on the rotation interval [1]. In any case, the knowledge of the rotation interval of circle maps of degree one is of theoretical importance.

The rotation number was introduced by H. Poincaré to study the movement of celestial bodies [14], and since then has been found to model a wide variety of physical and sociological processes. The application to voting theory [8] [12] is specially surprising in this context.

E-mail addresses: alseda@mat.uab.cat, llalseda@crm.cat, sborros@mat.uab.cat.

Date: December 6, 2020.

2020 Mathematics Subject Classification. Primary 37E10, Secondary 37E45.

Key words and phrases. Rotation number, circle maps, nondecreasing degree one lifting, algorithm.

Supported by the Spain’s “Agencial Estatal de Investigación” (AEI) grants MTM2017-86795-C3-1-P and MDM-2014-0445 within the “María de Maeztu” Program.
The computation of the rotation number for invertible maps of degree 1 from $S^1$ onto itself is well studied, and many very efficient algorithms exist for its computation [5, 13, 15, 16]. However, there is a lack of an efficient algorithm for the non-invertible and non-differentiable case.

In this article, we propose a method that allows us to compute the rotation interval for the non-invertible case. Our algorithm is based on the fact that we can compute exactly the rotation number of a natural subclass of the class of continuous non-decreasing degree one circle maps that have a constant section and a rational rotation number. From this algorithm we get an efficient way to compute exactly the rotation interval of a natural subclass of the continuous non-invertible degree one circle maps by using the so called upper and lower maps, which, when different, always have a constant section.

To check the efficiency of our algorithm will use it to compute some classical results such as a Devil’s Staircase. When doing so, we will compare the efficiency of our algorithm with the performance of some other algorithms that have been traditionally used under the hypothesis of non-invertibility. On the other hand, we will also compute the rotation interval and Arnold tongues for a variety of maps, in the same comparing spirit. These maps include the Standard Map and variants of it but have issues either with the differentiability, or even with the continuity. Of course these variants are not well suited for algorithms that strongly use differentiability.

The paper is organised as follows. In Section 2 the theoretical background will be set. In Section 3 the algorithm will be presented, and in Section 4 we will provide the mentioned examples of the use of the algorithm. Finally in Section 5 we will discuss the advantages and disadvantages of the proposed algorithm.

2. A short Survey on Rotation Theory and the Computation of Rotation Numbers

We will start by recalling some results from the rotation theory for circle maps. To do this we will follow [2].

The floor function (i.e. the function that returns the greatest integer less than or equal to the variable) will be denoted as \([\cdot]\). Also the decimal part of a real number $x \in \mathbb{R}$, defined as $x - [x] \in [0, 1)$ will be denoted by $\{x\}$.

In what follows $S^1$ denotes the circle, which is defined as the set of all complex numbers of modulus one. Let $e: \mathbb{R} \to S^1$ be the natural projection from $\mathbb{R}$ to $S^1$, which is defined by $e(x) := \exp(2\pi ix)$.

Let $f: S^1 \to S^1$ be continuous map. A continuous map $F: \mathbb{R} \to \mathbb{R}$ is a lifting of $f$ if and only if $e(F(x)) = f(e(x))$ for every $x \in \mathbb{R}$. Note that the lifting of a circle map is not unique, and that any two liftings $F$ and $F'$ of the same continuous map $f: S^1 \to S^1$ verify $F = F' + k$ for some $k \in \mathbb{Z}$.

For every continuous map $f: S^1 \to S^1$ there exists an integer $d$ such that

$$F(x + 1) = F(x) + d$$

for every lifting $F$ of $f$ and every $x \in \mathbb{R}$ (that is, the number $d$ is independent of the choice of the lifting and the point $x \in \mathbb{R}$). We shall call this number $d$ the degree of $f$. The degree of a map roughly corresponds to the number of times that the whole image of the map $f$ covers homotopically $S^1$.

In this paper we are interested studying maps of degree 1, since the rotation theory is well defined for the liftings of these maps.

We will denote the set of all liftings of maps of degree 1 by $L_1$. Observe that to define a map from $L_1$ it is enough to define $F|_{[0,1]}$ (see Figure 1) because $F$ can be globally defined as $F(x) = F|_{[0,1]}(\{x\}) + [x]$ for every $x \in \mathbb{R}$.
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Figure 1. An example of a map from $\mathcal{L}_1$ which can be considered as a toy model for the elements of that class. The picture shows $F|_{[0,1]}$ and $F$ is globally defined as $F(x) = F|_{[0,1]}(\{x\}) + \lfloor x \rfloor$.

Remark 2.1. It is easy to see that, for every $F \in \mathcal{L}_1$, $F^n(x + k) = F^n(x) + k$ for every $n \in \mathbb{N}$, $x \in \mathbb{R}$ and $k \in \mathbb{Z}$. Consequently, $F^n \in \mathcal{L}_1$ for every $n \in \mathbb{N}$.

Definition 2.2. Let $F \in \mathcal{L}_1$, and let $x \in \mathbb{R}$. We define the rotation number of $x$ as

$$\rho_F(x) := \limsup_{n \to \infty} \frac{F^n(x) - x}{n}.$$  

Observe (Remark 2.1) that, $\rho_F(x) = \rho_F(x + k)$ for every $k \in \mathbb{Z}$. The rotation set of $F$ is defined as:

$$\text{Rot}(F) = \{\rho_F(x) : x \in \mathbb{R}\} = \{\rho_F(x) : x \in [0,1]\}.$$  

Ito [6], proved that the rotation set is a closed interval of the real line. So, henceforth the set $\text{Rot}(F)$ will be called the rotation interval of $F$.

Proposition 2.3. Let $F \in \mathcal{L}_1$ be non-decreasing. Then, for every $x \in \mathbb{R}$ the limit

$$\lim_{n \to \infty} \frac{F^n(x) - x}{n}$$

exists and is independent of $x$.

For a non-decreasing map $F \in \mathcal{L}_1$, the number $\rho_F(x) = \lim_{n \to \infty} \frac{F^n(x) - x}{n}$ will be called the rotation number of $F$, and will be denoted by $\rho_F$.

Now, by using the notation from [2], we will introduce the notion of upper and lower functions, that will be crucial to compute the rotation interval.

Definition 2.4. Given $F \in \mathcal{L}_1$ we define the $F$-upper map $F_u$ as

$$F_u(x) := \sup\{F(y) : y \leq x\}.$$  

Similarly we will define the $F$-lower map as

$$F_l(x) := \inf\{F(y) : y \geq x\}.$$  

An example of such functions is shown in Figure 2.

It is easy to see that $F_l, F_u \in \mathcal{L}_1$ are non decreasing, and $F_l(x) \leq F(x) \leq F_u(x)$ for every $x \in \mathbb{R}$.
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Figure 2. An example of a map $F \in \mathcal{L}_1$ with its lower map $F_l$ in red and its upper map $F_u$ in blue.

The rationale behind introducing the upper and lower functions comes from the following result, stating that the rotation interval of a function $F \in \mathcal{L}_1$ is given by the rotation number of its upper and lower functions.

**Theorem 2.5.** Let $F \in \mathcal{L}_1$. Then,

$$\text{Rot}(F) = \left[ \rho_{F_l}, \rho_{F_u} \right].$$

Note that this theorem makes indeed sense, since the upper and lower functions are non-decreasing and by Proposition 2.3 they have a single well defined rotation number.

Let $f : S^1 \to S^1$ and let $z \in S^1$. The $f$-orbit of $z$ is defined to be the set

$$\text{Orb}_f(z) := \{ z, f(z), f^2(z), \ldots, f^n(z), \ldots \}.$$

We say that $z$ is an $n$-periodic point of $f$ if $\text{Orb}_f(z)$ has cardinality $n$. Note that this is equivalent to $f^n(z) = z$ and $f^k(z) \neq z$ for every $k < n$. In this case the set $\text{Orb}_f(z)$ will be called an $n$-periodic orbit (or, simply, a periodic orbit).

If we have a periodic orbit of a circle map, a natural question that might arise is how it behaves at a lifting level. This motivates the introduction of the notion of a lifted cycle.

Given a set $A \subset \mathbb{R}$ and $m \in \mathbb{Z}$ we will denote $A + m := \{ x + m : x \in A \}$. Analogously, we set

$$A + Z := \{ x + m : x \in A, m \in \mathbb{Z} \}.$$

**Definition 2.6.** Let $f : S^1 \to S^1$ be a continuous map and let $F$ be a lifting of $f$. A set $P \subset \mathbb{R}$ is called a lifted cycle of $F$ if $e(P)$ is a periodic orbit of $f$. Observe that, then $P = P + Z$. The period of a lifted cycle is, by definition, the period of $e(P)$. Hence, when $e(P)$ is an $n$-periodic orbit of $f$, $P$ is called an $n$-lifted cycle, and every point $x \in P$ will be called an $n$-periodic (mod 1) point of $F$.

The relation between lifted orbits and rotation numbers is clarified by the next lemma.

**Lemma 2.7.** Let $F \in \mathcal{L}_1$. Then, $x$ is an $n$-periodic (mod 1) point of $F$ if and only if there exists $k \in \mathbb{Z}$ such that $F^n(x) = x + k$ but $F^j(x) - x \notin \mathbb{Z}$ for $j = 1, 2, \ldots, n - 1$. In this case,

$$\rho_F(x) = \lim_{m \to \infty} \frac{F^m(x) - x}{m} = \frac{k}{n}.$$
Moreover, let $P$ be a lifted $n$-cycle of $F$. Every point $x \in P$ is an $n$-periodic (mod 1) point of $F$, and the above number $k$ does not depend on $x$. Hence, for every $x \in P$ we have $\rho_x(P) := \rho_x(x) = \frac{k}{n}$.

Now we can revisit Proposition 2.3.

**Proposition 2.3.** Let $F \in \mathcal{L}_1$ be non-decreasing. Then, for every $x \in \mathbb{R}$ the limit

$$\rho_x := \lim_{n \to \infty} \frac{F^n(x) - x}{n}$$

exists and is independent of $x$. Moreover, $\rho_x$ is rational if and only if $F$ has a lifted cycle.

In the next two subsections we will survey on two known algorithms that have already been used to compute rotation numbers of non-differentiable and non-invertible liftings from $\mathcal{L}_1$. The first one (Algorithm 1) stems automatically from the definition of rotation number (Definition 2.2); the other one (Algorithm 2) is due to Simó et al. [7].

### 2.1. Algorithm 1

**Algorithm 1**

**Algorithm pseudocode**

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Algorithm 1: the numerical algorithm to compute the rotation interval that stems from the definition of rotation number.

First algorithm to compute $\rho_x$ consists in using Proposition 2.3 and the following approximation, for $n$ large enough in relation to the desired tolerance:

$$\rho_x = \lim_{n \to \infty} \frac{F^n(x) - x}{n} \approx \frac{F^n(0)}{n}.$$

The implementation of the computation of this approximation to the rotation number can be found in the side algorithm pseudocode.

Since the maps from $\mathcal{L}_1$ are defined so that $F(x) = F|_{[0,1]}(\lfloor x \rfloor) + \lfloor x \rfloor$, we need to evaluate the function $\text{FLOOR}(\cdot) = \lfloor \cdot \rfloor$ once per iterate. So, for clarity and efficiency, it seems advisable to split $F^n(0)$ as $\lfloor F^n(0) \rfloor + [F^n(0)]$. The next lemma clarifies the computation error as a function of the number of iterates. In particular it explicitly gives the necessary number of iterates, given a fixed tolerance.

For every non-decreasing lifting $F \in \mathcal{L}_1$ and every $n \in \mathbb{N}$ we set (see Figure 3)

$$\ell_F(n) := \min_{x \in \mathbb{R}} |F^n(x) - x| = \min_{x \in [0,1]} |F^n(x) - x|.$$

The second equality holds because $F$ has degree 1, and hence $\ell_F(n)$ is well defined.

**Lemma 2.8.** For every non-decreasing lifting $F \in \mathcal{L}_1$ and $n \in \mathbb{N}$ we have

(a) either $F^n(z) = z + \ell_F(n) + 1$ for some $z \in \mathbb{R}$, or $x + \ell_F(n) \leq F^n(x) < x + \ell_F(n) + 1$ for every $x \in \mathbb{R}$;

(b) $\ell_F(n) \leq \rho_x \leq \ell_F(n) + 1$; and

(c) $|\rho_x - \frac{F^n(x) - x}{n}| < \frac{1}{n}$ for every $x \in \mathbb{R}$.
Proposition 2.3 and Lemma 2.7 imply that which proves (c) in this case. Now observe that from the definition of we will prove the whole lemma by considering two alternative cases. Assume first that Proposition 2.3 and Lemma 2.7 imply that \( \ell_F(n) \leq \lfloor F^n(x) - x \rfloor \leq F^n(x) - x \) for every \( x \in \mathbb{R} \). Moreover, it is implicitly assumed that \( \rho_r \) holds in this case. Now observe that from the definition of \( \ell_F(n) \) we have

\[
\ell_F(n) \leq \lfloor F^n(x) - x \rfloor \leq F^n(x) - x
\]

which proves (c) in this case.

Now we consider the case for every \( x \in \mathbb{R} \). In view of the definition of \( \ell_F(n) \), we cannot have

\[
F^n(x) - x > \ell_F(n) + 1
\]

for every \( x \in \mathbb{R} \). Hence, by the continuity of \( F^n(x) - x \) and (1),

\[
(2) \quad \ell_F(n) \leq F^n(x) - x < \ell_F(n) + 1
\]

for every \( x \in \mathbb{R} \). This proves (a).

Now we prove (b). We consider the functions: \( x \mapsto \ell_F(n) + x, \) \( F^n, \) and \( x \mapsto \ell_F(n) + 1 + x. \) They are all non-decreasing and, by Remark 2.1, they belong to \( L_1 \). Hence, by Proposition 2.3 and Lemma 3.7.19 and (2),

\[
\ell_F(n) = \rho_{x=\ell_F(n)} \leq \rho_{x=\ell_F(n)+1} = \ell_F(n) + 1.
\]

Consequently,

\[
\frac{\ell_F(n)}{n} \leq \rho_r = \frac{\rho_{x=n}}{n} \leq \frac{\ell_F(n)}{n} + 1,
\]

and (b) holds. Moreover, (2) is equivalent to

\[
\frac{\ell_F(n)}{n} \leq \frac{F^n(x) - x}{n} \leq \frac{\ell_F(n) + 1}{n},
\]

which proves (c).

\[\square\]

2.2. Algorithm 2: the Simó et al. algorithm to compute the rotation interval. First of all, it should be noted that even though the authors propose an algorithm to compute the rotation interval for a general map \( F \in L_1 \), we will only use it for non-decreasing maps. \textit{A priori} this algorithm is radically different from Algorithm 1 and it gives an estimate of \( \rho_r \) by providing and upper and a lower bound of the rotation number (rotation interval in the original paper) of \( F \). Moreover, it is implicitly assumed that \( \rho_r \in [0,1] \) (in particular that \( F(0) \in [0,1] \) — this can be achieved by replacing the lifting \( F \) by the lifting \( G := F - \lfloor F(0) \rfloor \), if necessary). The algorithm goes as follows:

\begin{figure}[h]
\centering
\includegraphics[width=0.5\textwidth]{algorithm2.png}
\caption{Plot of \( x + \ell_F(n) \) and \( x + \ell_F(n)+1 \), and \( F^n(x) \) for two arbitrary non-decreasing maps \( F \in L_1 \) that fit in the two cases of the lemma.}
\end{figure}
Algorithm 2 Simó et al. (7) Algorithm in pseudocode

procedure Rotation_Number(F, n)
    index[] ←
    x ← 0
    ρ_min ← 0
    ρ_max ← 1
    for i ← 0, n do
        x ← F(x)
        k_i ← floor(x)
        α_i ← x - k_i
        index[i] ← i
    end for
    sort α[index[i]] by rearranging index[]
    for i ← 0, n - 1 do
        ρ_aux ← \frac{k_{index[i+1]} - k_{index[i]}}{index[i+1] - index[i]}
        if index[i+1] > index[i] then
            ρ_min ← max{ρ_min, ρ_aux}
        else
            ρ_max ← min{ρ_max, ρ_aux}
        end if
    end for
    return ρ_min, ρ_max
end procedure

(Alg. 2-1) Decide the number of iterates n in function of a given tolerance.
(Alg. 2-2) For i = 0, 1, 2, ..., n compute k_i = \lfloor F^i(x_0) \rfloor and α_i = F^i(x_0) - k_i (i.e.
α_i is the fractionary part of F^i(x_0)).
(Alg. 2-3) Sort the values of k_i and α_i so that α_0 < α_1 < ... < α_n (this can be achieved efficiently with the help of an index vector).
(Alg. 2-4) Initialise ρ_min = 0 and ρ_max = 1.
(Alg. 2-5) For j = 0, 1, 2, ..., n - 1 set ρ_j = \frac{k_{i_j+1} - k_{i_j}}{i_{j+1} - i_j}, and
    • if i_{j+1} > i_j set ρ_min = max{ρ_min, ρ_j}; otherwise,
    • if i_{j+1} < i_j set ρ_max = min{ρ_max, ρ_j}.
(Alg. 2-6) Return ρ_max and ρ_min as upper and lower bounds of the rotation number of F, respectively.

The real issue in this algorithm consists in dealing with the error. If the rotation number ρ_F satisfies a Diophantine condition \left| ρ_F - \frac{p}{q} \right| ≤ cq^{-ν}, with c > 0 and ν ≥ 2,
then the error verifies
\[ ε < \frac{1}{(cn^ν)^\frac{1}{ν+1}}. \]

Note that this error depends strongly on the chosen number n of iterates, and that n must be chosen before knowing what the rotation number could possibly be. Hence Algorithm 2 is not well suited to compute unknown rotation numbers of \textit{L}_1 maps. However, it is excellent in continuation methods where the current rotation number gives a good estimate of the next one.

Remark 2.9. Note that the original aim of the algorithm to determine the existence of closed invariant curves on dynamical systems on the plane rather than the computation of rotation numbers of a given map of the circle. The rationale of the algorithm is that if, after computing ρ_min and ρ_max, we find that ρ_min > ρ_max then the computed orbit cannot lay on a closed invariant curve. This explains most of
the limitations we have encountered, such as the lack of an \textit{a priori} estimate of the error, or the fact that the algorithm is suited only for rotation numbers \( \rho \in [0, 1] \).

3. An algorithm to compute rotation numbers of non-decreasing maps with a constant section

The \textit{diameter of an interval} \( K \) which, by definition is equal to the absolute value of the difference between their endpoints, will be denoted as \( \text{diam}(K) \).

A \textit{constant section} of a lifting \( F \) of a circle map is a closed non-degenerate (i.e. different from a point or, equivalently, with non-empty interior, or such that \( \text{diam}(K) > 0 \)) subinterval of \( \mathbb{R} \) such that \( F|_K \) is constant. In the special case when \( F \in \mathcal{L}_1 \), we have that \( F(x + 1) = F(x) + 1 \neq F(x) \) for every \( x \in \mathbb{R} \). Hence, \( \text{diam}(K) < 1 \).

The algorithm we propose is based on Lemma 2.8 but, specially, on the following simple proposition which allows us to compute \textit{exactly} the rotation number of a non-decreasing lifting from \( \mathcal{L}_1 \) that has a constant section, provided that \( F^n(K) \cap (K + \mathbb{Z}) \neq \emptyset \). In this sense, Proposition 3.1 has a completely different strategical aim than Algorithm 1 and Lemma 2.8 which try to (costly) estimate the rotation number.

**Proposition 3.1.** Let \( F \in \mathcal{L}_1 \) be non-decreasing and have a constant section \( K \). Assume that there exists \( n \in \mathbb{N} \) such that \( F^n(K) \cap (K + \mathbb{Z}) \neq \emptyset \), and that \( n \) is minimal with this property. Then, there exists \( \xi \in \mathbb{R} \) such that \( F^n(K) = \{ \xi \} \subset K + m \) with \( m = \lfloor \xi - \min K \rfloor \in \mathbb{Z} \), \( \xi \) is an \( n \)-periodic (mod 1) point of \( F \), and \( \rho_F = \frac{m}{n} \).

**Proof.** Since \( K \) is a constant section of \( F \), \( F(K) \) contains a unique point, and hence there exists \( \xi \in \mathbb{R} \) such that \( F^n(K) = \{ \xi \} \). Then, the fact that \( F^n(K) \cap (K + \mathbb{Z}) \neq \emptyset \) implies that \( \xi \in K + m \) with \( m = \lfloor \xi - \min K \rfloor \in \mathbb{Z} \).

Set \( \tilde{\xi} := \xi - m \in K \). Then, \( \{ F^n(\tilde{\xi}) \} = F^n(K) = \{ \tilde{\xi} + m \} \). Moreover, the minimality of \( n \) implies that \( F^j(\tilde{\xi}) - \tilde{\xi} \notin \mathbb{Z} \) for \( j = 1, 2, \ldots, n - 1 \). So, Lemma 2.7 tells us that \( \tilde{\xi} \) (and hence \( \xi \)) is an \( n \)-periodic (mod 1) point of \( F \). Thus, \( \rho_F = \frac{m}{n} \) by Proposition 2.3.

As already said, Proposition 3.1 is a tool to compute \textit{exactly} the rotation numbers of non-decreasing liftings \( F \in \mathcal{L}_1 \) which have a constant section and have a lifted cycle intersecting the constant section (and hence having rational rotation number). In the next subsection we shall investigate how restrictive are these conditions, when dealing with computation of rotation numbers.

3.1. On the genericity of Proposition 3.1

First observe that the fact that Proposition 3.1 only allows the computation of rotation numbers of non-decreasing liftings \( F \in \mathcal{L}_1 \) which have a constant section is not restrictive at all. Indeed, if we want to compute rotation intervals of \textit{non-invertible} continuous circle maps of degree one, Theorem 2.5 tells us that this is exactly what we want.

Clearly, one of the real restrictions that cannot be overcome in the above method to compute \textit{exact} rotation numbers is that it only works for maps having a rational rotation number.

On the other hand, we also have the formal restriction that Proposition 3.1 requires that the map \( F \) has a lifted cycle intersecting the constant section (indeed this is a consequence of the condition \( F^n(K) \cap (K + \mathbb{Z}) \neq \emptyset \)). A natural question is whether this restriction is just formal or it is a real one. In the next example we will see that the restriction is not superfluous since there exist maps which do not satisfy it.

Consequently, Proposition 3.1 is useless in computing the rotation numbers of non-decreasing liftings in \( \mathcal{L}_1 \) which have a constant section and either irrational
rotation number or rational rotation number but do not have any lifted cycle intersecting the constant section. The only reasonable solution to these problems is to use an iterative algorithm to estimate the rotation number with a prescribed error, such as Algorithm 1, Algorithm 2 or others.

Example 3.2. There exist non-decreasing liftings in $\mathcal{L}_1$ which have a constant section and rational rotation number but do not have any lifted cycle intersecting the constant section. Let $F \in \mathcal{L}_1$ be the map such that $F(x) = F_{[0,1]}(x) + \lfloor x \rfloor$ for every $x \in \mathbb{R}$, and let

$$F_{[0,1]}(x) := \begin{cases} 
    x + 0.2 & \text{if } x \in [0,0.1], \\
    \frac{x}{2} + 0.25 & \text{if } x \in [0.1,0.3], \\
    7x - 1.7 & \text{if } x \in [0.3,0.4], \\
    \frac{x}{4} + 1 & \text{if } x \in [0.4,0.8], \\
    1.2 & \text{if } x \in [0.8,1].
\end{cases}$$

Figure 4. Example of a non-decreasing lifting in $\mathcal{L}_1$ with a constant section and rational rotation number which does not verify the assumptions of Proposition 3.1. The map $F$ is a non-decreasing lifting from $\mathcal{L}_1$, having a constant section $K = [0.8,1]$ and rotation number $\frac{1}{3}$ given by the 3-lifted cycle $P = \{0.1,0.3,0.4\} + \mathbb{Z}$ (c.f. Lemma 2.7 and Proposition 2.3).

Now let us see that $F$ does not have any lifted cycle intersecting the constant section. First, observe that $F^3(K) = F(F(F(K))) = F(F(\{1.2\})) = F(\{1.35\}) = \{1.75\} \not\subset K + \mathbb{Z}$.

Hence, there is no lifted cycle of period 3 intersecting $K$. On the other hand, again by Lemma 2.7, we have that if $x$ is an $n$-periodic (mod 1) point of $F$ then there exists $k \in \mathbb{Z}$ such that $F^n(x) = x + k$ and

$$\frac{1}{3} = \rho_F = \lim_{m \to \infty} \frac{F^m(x) - x}{m} = \rho_F(x) = \frac{k}{n}.$$

Moreover, since $F$ is non-decreasing, we know by [2, Corollary 3.7.6] that $n$ and $k$ must be relatively prime. Thus, any lifted cycle of $F$ has period 3, and from above this implies that there is no lifted cycle intersecting $K$.

3.2. Algorithm 3. A constant section based algorithm arising from Proposition 3.1. From the last paragraph of the previous subsection it becomes evident that Proposition 3.1 does not give a complete algorithm to compute rotation numbers of non-decreasing liftings in $\mathcal{L}_1$ which have a constant section. Such an algorithm must rather be a mix-up of Proposition 3.1 and Algorithm 1 to be used when we are not able to determine whether we are in the assumptions of that proposition. As in Algorithm 1, for efficiency and because Proposition 3.1 requires the computation of $m$ as an integer part, we will split $F^n(0)$ as $\lfloor F^n(0) \rfloor + [F^n(0)]$ (here we are denoting the constant section by $K$ and assuming that $0 \in K$ — to be justified later). Then, observe that the computations to be performed are exactly...
Algorithm 3
Constant Section Based Algorithm

For a non-decreasing map $F \in \mathcal{L}_1$ parametrised so that $[-\text{tol}, \beta + \text{tol}]$ is a constant section of $F$

**define** \( \text{tol} \leftarrow \) \( \triangledown \) Procedure parameter that bounds the rounding errors in the computation of $F^n(0)$

**procedure** Rotation_Number\((F, \beta, \text{error})\)

\( \text{max\_iter} \leftarrow \text{cei} \left( \frac{1}{\text{ERROR}} \right) \) \( \triangledown \) Maximum number of iterates allowed (to estimate the rotation number with the prescribed error when reached)

\( x \leftarrow 0 \)

\( m \leftarrow 0 \)

for \( n \leftarrow 1, \text{max\_iter} \) do

\( x \leftarrow F(x) \)

\( s \leftarrow \text{FLOOR}(x) \)

\( m \leftarrow m + s \)

\( x \leftarrow x - s \) \( \triangledown x = \llbracket F^n(0) \rrbracket = F^n(0) - m \)

if \( x \leq \beta \) then

\( \text{return} \ \frac{m}{n} \) \( \triangledown \) Exact rotation number: Proposition 3.1 holds assuming that the rounding error of $F^n(0)$ is smaller than \( \text{tol} \)

end if

end for

\( \text{return} \ \frac{m + x}{\text{max\_iter}} \) \( \triangledown \) We do not know whether we are in the assumptions of Proposition 3.1. So, we iteratively estimate the rotation number as in Algorithm 1.

The error bound is given by Lemma 2.8.

the same in both cases (meaning when we can use Proposition 3.1 and when alternatively we must end up by using Algorithm 1; except for the conditionals that check whether there exists \( n \leq \text{max\_iter} \) such that $F^n(K) \cap (K + \mathbb{Z}) \neq \emptyset$ is verified (that is, whether the assumptions of Proposition 3.1 are verified) before exhausting the max_iter iterates determined a priori.

In what follows $\hat{F^n}(0)$ will denote the computed value of $F^n(0)$ with rounding errors for $n = 1, 2, \ldots, \text{max\_iter}$. The algorithm goes as follows (see Algorithm 3 for a full implementation in pseudocode, and see the explanatory comments below):

(Alg. 3-1) Decide the maximum number of iterates max_iter = cei \left( \frac{1}{\text{ERROR}} \right) to perform in the worst case (i.e. when Proposition 3.1 does not work).

(Alg. 3-2) Re-parametrize the lifting $F$ so that it has a maximal (with respect to the inclusion relation) constant section of the form $[-\text{tol}, \beta + \text{tol}]$, where \( \text{tol} \) is the pre-defined rounding error bound.

(Alg. 3-3) Initialize $x = 0$ and $m = 0$.

(Alg. 3-4) Compute iteratively $x = \llbracket \hat{F^n}(0) \rrbracket$ and $m = \lfloor \hat{F^n}(0) \rfloor$ (so that $\hat{F^n}(0) = x + m$) for $n \leq \text{max\_iter}$.

(Alg. 3-5) Check whether $x \leq \beta$. On the affirmative we are in the assumptions of Proposition 3.1 and thus, $\rho_F = \frac{m}{n}$. Then, the algorithm returns this value as the “exact” rotation number.

(Alg. 3-6) If we reach the maximum number of iterates (i.e. $n = \text{max\_iter}$) without being in the assumptions of Proposition 3.1 (i.e. with $x \gt \beta$ for every $x$) then, by Lemma 2.8, we have

$$|\rho_F - \frac{m + x}{\text{max\_iter}}| = |\rho_F - \frac{\hat{F^n}(0)}{\text{max\_iter}}| \approx |\rho_F - \frac{F^n(0)}{\text{max\_iter}}| < \frac{1}{\text{max\_iter}},$$
and the algorithm returns \( \frac{m+x}{\max_{\text{iter}}} \) as an estimate of \( \rho_F \) with \( \frac{1}{\max_{\text{iter}}} \) as the estimated error bound.

**Remark 3.3.** The fact that we can only check whether the assumptions of Proposition 3.1 are verified before exhausting the \( \max_{\text{iter}} = \text{ceil}(\frac{1}{\text{error}}) \) iterates determined a priori does not allow to take into account that \( F \) may have a lifted cycle intersecting the constant section but of very large period, i.e. with period larger than \( \max_{\text{iter}} \). In practice this problem is totally equivalent to the non-existence (or rather invisibility) of a lifted cycle intersecting the constant section, and it can be considered as a new (algorithmic) restriction to Proposition 3.1. It is solved in (Alg. 3-6) in the same manner as the two other problems related with the applicability of Proposition 3.1 that have already been discussed: by estimating the rotation number as in Algorithm 1.

In the last part of this subsection we are going to discuss the rationale of (Alg. 3-2) (and, as a consequence of (Alg. 3-5)). The necessity of this tuning of the algorithm comes again from a challenge concerning the application of Proposition 3.1, which turns to be one of the most relevant restrictions in the use of that proposition. We will begin by discussing how we can efficiently check the condition \( \xi = F^n(0) \in K + \mathbb{Z} \) (or equivalently \( F^n(K) \cap (K + \mathbb{Z}) \neq \emptyset \)) by taking into account that the computation of \( F(x) \) is done with rounding errors, and thus we do not know the exact values of \( F^n(0) \) for \( n = 1, 2, \ldots, \max_{\text{iter}} \). The next example shows the problems arising in this situation.

**Example 3.4.** \( F^n(0) \in K + \mathbb{Z} \) but \( F^n(K) \cap (K + \mathbb{Z}) = \emptyset \), and this leads to a completely wrong estimate of \( \rho_F \).

Let \( F \in \mathcal{L}_1 \) be the map such that \( F(x) = F\mid_{[0,1]}(\{x\}) + \lfloor x \rfloor \) for every \( x \in \mathbb{R} \), and let

\[
F\mid_{[0,1]}(x) := \begin{cases} 
\frac{4}{3}x + \mu & \text{if } x \in [0, \frac{3}{4}], \\
1 + \mu & \text{if } x \in [\frac{3}{4}, 1],
\end{cases}
\]

with \( \mu = \frac{819}{3124} - 10^{-16} \).

For this map \( F \) we have \( K = [-\frac{3}{4}, 0] \) and (see Figure 3) the graph of \( F^5 \) lies above the graph of \( x \mapsto x + 1 \) and below the graph of \( x \mapsto x + 2 \), but very close to it at five \( F \)-preimages of \( x = \frac{3}{4} \). On the other hand,

\[
F^5(0) = 1.7499999999999887 \cdots \notin K + \mathbb{Z}
\]

but the distance between \( F^5(0) \) and \( K + \mathbb{Z} \) is \( \frac{7}{4} - F^5(0) \approx 1.138 \cdot 10^{-15} \). Should the computations be done with rounding errors of this last magnitude, we may have \( \tilde{F}^5(0) \approx \frac{7}{4} \) and accept erroneously that \( F^5(0) \in K + \mathbb{Z} \). This would lead to the conclusion that \( \rho_F = \frac{7}{4} \) but, as it can be checked numerically, \( \rho_F \approx 0.3983 \) which is far from \( \frac{7}{4} \).

At a first glance this seems to be paradoxical but, indeed, it can be viewed in the following way: The graph of \( F^5 \) does not intersect the diagonal (modulo 1)
x + 2, but there is a map $G$ close (at rounding errors distance) to $F$ such that the graph of $G^5$ intersects that diagonal, and this gives a lifted periodic orbit of period 5 and rotation number $\frac{2}{5}$ for $G$. On the other hand, nothing is granted about the modulus of continuity of $\rho_F$ as a function of $F$ (notice that that everything here is continuous including the dependence of the rotation number of $F$ on the parameter $\mu$), and this example explicitly shows that it may be indeed very big. In short, close functions can have very different rotation numbers.

The most reasonable solution to the problem pointed out in the previous example consists in restricting the size of $K$ depending of an a priori estimate of the rounding errors in computing $\tilde{F}_n(0)$ for $n = 1, 2, \ldots, \max_{\text{iter}}$. Thus, we denote by $\text{tol}$ an upper bound of these rounding errors, so that
\[
|F_n(0) - \tilde{F}_n(0)| \leq \text{tol}
\]
holds for $n = 1, 2, \ldots, \max_{\text{iter}},$

and, given a maximal (with respect to the inclusion relation) constant section $K$ such that $0 \in K$ we write $K := [\alpha - \text{tol}, \beta + \text{tol}]$. Then observe that the condition $\tilde{F}_n(0) \in [\alpha, \beta] + m$ for some $n \in \mathbb{N}$ and $m \in \mathbb{Z}$ implies $\xi = F^n(0) \in K + m$, and $\rho_F = \frac{m}{n}$ by Proposition 3.1.

In practice, this “rounding errors free” version of the algorithm imposes a new restriction to the applicability of Proposition 3.1 (in the sense that it reduces even more the class of functions for which we can get the “exact rotation number”). However, as before, the rotation numbers of the maps in the assumptions of Proposition 3.1 for which we cannot compute the “exact rotation number” can be estimated as in Algorithm 1.

The computational efficiency of the algorithm strongly depends on how we check the condition $\tilde{F}_n(0) \in K + \mathbb{Z}$. Taking into account the above considerations and improvements of the algorithm, this amounts checking whether $\alpha + \ell \leq \tilde{F}_n(0) \leq \beta + \ell$ for some $\ell \in \mathbb{Z}$, and we have to do so by using $x = \lfloor \tilde{F}_n(0) \rfloor$ and $m = \lfloor \tilde{F}_n(0) \rfloor$ instead of $\tilde{F}_n(0) = x + m$, which is the algorithmic available information. Checking whether $\alpha + \ell \leq \tilde{F}_n(0) \leq \beta + \ell$ for some $\ell \in \mathbb{Z}$ is problematic since it requires

\begin{figure}[h]
\centering
\includegraphics[width=\textwidth]{graph.png}
\caption{The graph of $F^5$. It lies below the graph of $x \mapsto x + 2$ but very close to it at five $F$-preimages of $x = \frac{3}{4}$.}
\end{figure}
at least two comparisons, and moreover in general $\ell \neq m$ (and thus we need some more computational effort to find the right value of $\ell$). A very easy solution to this problem is to change the parametrization of $F$ so that $\alpha = 0$. In this situation we have

$$m = m + \alpha \leq \widehat{F^n}(0), \quad m + \beta < m + 1$$

because $\text{diam}(K) < 1$, and $m = \left\lceil \widehat{F^n}(0) \right\rceil$. Consequently, $\alpha + \ell \leq \widehat{F^n}(0) \leq \beta + \ell$ for some $\ell \in \mathbb{Z}$ is equivalent to

$$\ell = m \quad \text{and} \quad x \leq \beta.$$ 

Thus, by “tuning” $F$ so that $\alpha = 0$ we get that $\ell = m$ and we manage to determine whether $\widehat{F^n}(0) \in [\alpha, \beta] + m$ just with one comparison ($x \leq \beta$).

To see that and how we can change the parametrization of $F$ (that is the point 0) so that $\alpha = 0$ consider the map $G(x) := F(x + \alpha) - \alpha$. Clearly, $F$ and $G$ are conjugate by the rotation of angle $\alpha$: $x \mapsto x + \alpha$. Then, obviously, $G$ is a non-decreasing map in $L_1$, has a constant section $[-\text{tol}, \beta - \alpha + \text{tol}]$, and $\rho_F = \rho_G$. So, every lifting can be replaced by one of its re-parametrizations with the same rotation number and constant section $[-\text{tol}, \beta + \text{tol}]$, where $\beta < 1 - 2\text{tol}$.

4. Testing the Algorithm

In this section we will test the performance of Algorithm 3 by comparing it against Algorithms 1 and 2 when dealing with different usual computations concerning rotation intervals. First we will compare the efficiency of the three algorithms in computing and plotting Devil’s Staircases. Afterwards we will plot rotation intervals and Arnold tongues for two bi-parametric families that mimic the standard map family. In the latter two cases, we will try to compare our algorithm with Algorithms 1 and 2 whenever possible.

4.1. Computing Devil’s staircases. In this subsection we will perform the comparison of algorithms by computing and plotting the Devil’s staircase for the para-

metric family $\{F_\mu\}_{\mu \in [0,1]} \subset L_1$ defined as

**Definition 4.1.**

$$F_\mu(x) = F_\mu\big|_{[0,1]}(\langle x \rangle) + \lfloor x \rfloor,$$

where (see Figure 1)

$$F_\mu\big|_{[0,1]}(x) = \begin{cases} 
\frac{4}{3}x + \mu & \text{if } x \leq \frac{3}{4} \\
\mu + 1 & \text{if } x > \frac{3}{4}.
\end{cases}$$

Before doing this we shall remind the notion of a Devil’s Staircase, and why typically exist for such families. To this end we will first recall and survey on the notion of persistence of a rotation interval.

**Definition 4.2.** Given a subclass $A$ of $L_1$, we say that $F \in A$ has an $A$-persistent rotation interval if there exists a neighbourhood $U$ of $F$ in $A$ such that

$$\text{Rot}(G) = \text{Rot}(F)$$

for every $G \in U$.

We can now state the Persistence Theorem (c.f. [10]):

**Theorem 4.3** (Persistence Theorem). Let $A$ be a subclass of $L_1$. Then the following statements hold:

(a) The set of all maps with $A$-persistent rotation interval is open and dense in $A$ (in the topology of $A$).
(b) If $F$ has an $A$-persistent rotation interval, then $\rho_{F_1}$ and $\rho_{F_0}$ are rational.

**Remark 4.4.** If we apply Theorem 4.3 to our family $\{F_{\mu}\}_{\mu \in [0,1]}$ which verifies that the rotation number of $F_{\mu}$ exists for every $\mu \in [0,1]$, we have that the set of parameters $\mu \in [0,1]$ for which we have irrational rotation number has measure 0. Furthermore, for any $\kappa \in \mathbb{Q}$ such that there exists $\mu$ with $\rho_{F_\mu} = \kappa$, there exists an interval $[\alpha, \beta] \ni \mu$ such that for all $\eta \in [\alpha, \beta]$, $\rho_{F_\eta} = \kappa$.

The so-called Devil’s staircase is the result of plotting the rotation number as a function of the parameter $\mu$. By Theorem 4.3 we have that this plot will have constant sections for any rational rotation number, hence the ”Staircase” in the name.

To test the algorithms, a $\mu$-parametric grid computation of $\rho_{F_{\mu}}$ with $\mu$ ranging from 0 to 1 with a step of $10^{-5}$ has been done. For Algorithms 1 and 3 the error has been set to $10^{-6}$. For Algorithm 3 the tolerance has been set to $10^{-10}$. For Algorithm 2 we have arbitrarily set the number of iterates to 1000.

In Figure 6 we show a plot of the Devil’s Staircase computed with Algorithm 3 and the plots of the differences between $\rho_{F_{\mu}}$ computed with Algorithms 3 and 1, and the differences between $\rho_{F_{\mu}}$ computed with Algorithms 3 and 2.

Table 1 shows the times taken by each of the three algorithms in computing the

![Figure 6](image-url)
Table 1. Performance of the three algorithms studied for a variety of problems. The cells marked with N/A in blue remark that Algorithm 2 does not work in general for \( \rho \notin [0,1] \). The ones marked with N/A in red denote that the computation lasted more than a 100 processor hours and thus was terminated before it ended.

| Problem          | Function Family | Time taken by algorithm (s) |
|------------------|-----------------|-----------------------------|
| Devil’s Staircase| \( F_\mu \) (Def. 4.1) | 2425.25 210.648 0.1413 |
| Rotation Interval | PWLSM (Def. 4.7)   | 354.868 N/A 3.2874 |
|                  | DSM (Def. 4.8)    | 110.892 N/A 0.4737 |
|                  | DSM (Def. 4.7)    | 63.588 N/A 0.2463 |
| Arnol’d Tongues  | PWLSM            | N/A N/A 9729.17 |
|                  | DSM              | N/A N/A 4562.75 |

We remark that in the computation of the Devil’s Staircase, Algorithm 3 has been reduced to Algorithm 1 only for \( \mu = 0 \) and for \( \mu = 1 \), as one would expect, since these cases follow the pattern of Example 3.4.

As a part of the testing of the algorithms we have also considered the inverse problem: Given a value \( x \in \mathbb{R} \setminus \mathbb{Q} \) and a tolerance \( \varepsilon > 0 \) find the value \( \mu = \mu(x) \) such that \( \rho_{F_\mu} \in [x - \varepsilon, x + \varepsilon] \). This problem has turned to be extremely ill-conditioned: by choosing \( x \) to be an irrational such as the golden mean or \( \pi/4 \), the continuity module of the function \( \mu \mapsto \rho_{F_\mu} \) around \( \mu(x) \) was estimated to be at least \( 10^{25} \), making any attempt to solve the problem numerically a fool’s errand.

4.2. Rotation intervals for standard-like maps. In this subsection we test our algorithm by efficiently computing the rotation intervals and some Arnol’d tongues for three bi-parametric families of maps: the standard map family and two piecewise-linear extensions of it; one continuous but not differentiable, and another one which is not even continuous.

We emphasize that the usual algorithms such as the ones from [4, 13, 15, 16] cannot be used for these last two families families while the one we propose here it works like a charm.

First we will recall the notion of Arnol’d tongue.

**Definition 4.5 (Arnol’d Tongue [3]).** Let \( \{ F_{a,b} \}_{(a,b) \in P} \) be a two-parameter family of maps in \( \mathcal{L}_1 \) for which the rotation interval \( \text{Rot}(F_{a,b}) \) is well defined for every possible point \( (a,b) \in P \) in the parameter set. Given a point \( \varrho \in \mathbb{R} \) we define the \( \varrho \)-Arnold Tongue of \( \{ F_{a,b} \}_{(a,b) \in P} \) as

\[
\mathcal{T}_\varrho = \{(a,b) \in P : \varrho \in \text{Rot}(F_{a,b}) \} \subset P.
\]

Next we introduce each of the three families that we study and, for each of them we show the results and we explain the performance of the algorithm.

**Definition 4.6 (Standard Map).** \( S_{\Omega,a} \in \mathcal{L}_1 \) is defined as (see Figure 7):

\[
S_{\Omega,a}(x) := x + \Omega - \frac{a}{2\pi} \sin(2\pi x).
\]

\(^1\)The simulations have been done with an Intel® Core™ i7-3770 CPU @3.4GHz.
To compute the rotation intervals of $S_{\Omega,a}$ we will use Theorem 2.5 together with Algorithm 3. To this end, first we will compute $(S_{\Omega,a})_l$ and $(S_{\Omega,a})_u$ (that is, the lower and upper maps of $S_{\Omega,a}$), and then we will use Algorithm 3 to compute the rotation numbers $\rho(S_{\Omega,a})_l$ and $\rho(S_{\Omega,a})_u$ of these maps.

Note that $S_{\Omega,a}$ is non-invertible for $a > 1$. Hence, in this case, $(S_{\Omega,a})_l$ and $(S_{\Omega,a})_u$ do not coincide and have constant sections. However, the characterization of these constants sections is not straightforward, since their endpoints have to be computed numerically. This is the reason why the computations of the rotation intervals and Arnold’s tongues for the standard map have been the slowest ones.

In Figure 8 we show some graphs of the rotation interval and Arnold’s tongues for the Standard Map. The graphs of the rotation intervals are plotted for three different values of $\Omega$ as a function of the parameter $a$.

**Definition 4.7** (Piecewise-linear standard map). We start by defining a convenience map $\tau: [0,1] \to [-1,1]$ as follows:

\[
\tau(x) = \begin{cases} 
4x & \text{when } x \in [0, \frac{1}{4}], \\
2 - 4x & \text{when } x \in \left[\frac{1}{4}, \frac{3}{4}\right], \\
4(x-1) & \text{when } x \in \left[\frac{3}{4}, 1\right].
\end{cases}
\]

Then, the piecewise-linear standard map $T_{\Omega,a} \in L_1$ is defined by (see Figure 9):

\[
T_{\Omega,a}(x) = x + \Omega - \frac{a}{2\pi} \tau(\|x\|),
\]

which corresponds to the standard map but using the $\tau$ wave function instead of the $\sin(2\pi x)$ function.
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Figure 8. Graphs of the rotation interval and Arnol’d tongues for the Standard Map $S_{\Omega,a}$. The graphs of the rotation intervals are plotted as a function of the parameter $a$.

The upper and lower maps for this family are very easy to compute. Moreover, $T_{\Omega,a}$ is non-increasing for $a > \pi$ and hence, in this case, the upper and lower maps do not coincide and have constant sections.

Figure 9. The piecewise-linear standard map $T_{\Omega,a}$ with $a = \frac{5\pi}{2}$ and $\Omega = 0$. The lower map of $T_{\Omega,a}$ is drawn in blue, and the upper map in red.
To compute the rotation intervals and Arnol’d Tongues of $T_{\Omega,a}$ we proceed as for the Standard Map by using Theorem 2.5 and Algorithm 3.

In Figure 10 we show some graphs of the rotation interval and Arnol’d tongues for the piecewise-linear standard map. The graphs of the rotation intervals are plotted for three different values of $\Omega$ as a function of the parameter $a$.

![Graphs of rotation intervals and Arnol’d tongues for the piecewise-linear standard map $T_{\Omega,a}$.](image)

**Figure 10.** Graphs of the rotation interval and Arnol’d tongues for the piecewise-linear standard map $T_{\Omega,a}$. The graphs of the rotation intervals are plotted as a function of the parameter $a$.

**Definition 4.8 (The Discontinuous Standard Map).** $D_{\Omega,a} \in \mathcal{L}_1$ is defined as (see Figure 11):

\[
D_{\Omega,a}(x) := x + \Omega + \frac{a}{2\pi} \{x\}.
\]

The map $D_{\Omega,a}$, being discontinuous, belongs to the so called class of old heavy maps (the old part of the name stands for degree one lifting — that is, $D_{\Omega,a} \in \mathcal{L}_1$). A map $F \in \mathcal{L}_1$ is called heavy if for any $x \in \mathbb{R}$,

\[
\lim_{y \searrow x^+} F(y) \leq F(x) \leq \lim_{y \nearrow x^-} F(y)
\]

(in other words, the map “falls down” at all discontinuities).

Observe that for the class of old heavy maps the upper and lower maps in the sense of Definition 2.4 are well defined and continuous. Moreover, the whole family of water functions (c.f. [2]) is well defined and continuous. So, the rotation interval of the old heavy maps is well defined [11, Theorem A] and, moreover, Theorem 2.5 together with Algorithm 3 work for this class. Hence, to compute the rotation intervals and Arnol’d Tongues of $D_{\Omega,a}$ we proceed again as for the Standard Map.

As for the piecewise-linear standard maps the upper and lower maps are very easy to compute, and have constant sections for $a \neq 0$.

In Figure 12 we show some graphs of the rotation interval and Arnol’d tongues for the discontinuous standard map. The graphs of the rotation intervals are plotted...
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Figure 11. The discontinuous standard map with $a = 2\pi$ and $\Omega = 0$ with its lower map in blue and its upper map in red.

Figure 12. Graphs of the rotation interval and Arnol’d tongues for the discontinuous standard map $D_{\Omega,a}$. The graphs of the rotation intervals are plotted as a function of the parameter $a$.

for three different values of $\Omega$ as a function of the parameter $a$. The times taken for all the computation related with the rotation intervals and the Arnol’d Tongues for each of the families studied using Algorithms 1, 2 and 3 can be found in Table 1.
5. Conclusions

The algorithm proposed clearly outperforms all the other tested algorithms, both in precision and speed even though the “exact” (and quick) part of the algorithm does not work for all the non-decreasing liftings in $L_1$ which have a constant section (and hence the rotation number of these “bad” cases has to be computed with the much more inefficient classical algorithm). For all natural examples for which it has been tested, the computational speed and precision were unparalleled. Moreover, the set of functions becomes very general when one considers the fact that the upper and lower functions inherently have constant sections for any $F$ that is not strictly increasing. Hence, the algorithm becomes a crucial tool to compute rotation intervals for general functions in $L_1$ and hence to find the set of periods of such maps [2]. Moreover, a deeper study has been done on the dependence of the rotation number on the parameters. Our preliminary results have found that for irrational rotation numbers, the dependence of the parameters around them is extremely sensitive, with continuity module being at least $10^{25}$. This agrees with Theorem 4.3, which says that non-persistent functions have measure zero.

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