Comparison by simulation of PEGASIS and IEEPB routing protocols

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ABSTRACT

As the applications of wireless sensor networks (WSNs) became widely used throughout the years the importance of advanced sensor networks techniques increased as well. One of the main techniques used in WSNs is hierarchical routing which mainly aims to reduce the consumption of sensor nodes energy by assigning different roles to the sensor nodes to create multi-layer scheme for data transmission. This paper embraces a simulation for two known hierarchical routing protocols: Power-Efficient Gathering in Sensor Information Systems (PEGASIS) protocol and an Improved Energy-Efficient PEGASIS-Based (IEEPB) protocol. Both protocols aim to reduce the transmission distance in order to save the nodes energy by performing chain-based clustering. For evaluation, we measured the residual energy and control overhead throughout the network operation time and the results showed major flaws in both protocols such as long link problem and poor leader selection method in PEGASIS. Moreover, high nodes density problem in IEEPB.

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1. INTRODUCTION

A WSN is a network consist of large number of sensor nodes, i.e. tiny and low-cost devices that responsible of collecting data from the deployment area such as temperature, movement, light, sound, ... etc. Moreover, these nodes should collaborate to send their data to more powerful node known as the sink or base station (BS) in order to perform the appropriate actions. However, these kind of networks usually used in tough environments thus it is hard to recharge the sensor nodes [1]. Therefore, one of the principal challenges in WSNs is to reduce the energy consumption by sensors to ensure longer operation time of the networks.

In hierarchical routing protocols (HRPs), the nodes do not send directly to the BS; instead, some nodes are elected to act as higher-level nodes and receive data from the rest of the nodes then send the received data to the BS. Consequently, the transmission distance extremely reduced and the sensor nodes' life is prolonged. There are many routing protocols developed to improve the hierarchical routing process by following different approaches in organizing the roles among the sensor nodes and the nodes' connecting mechanism to set up the routing path. For instance, LEACH protocol [2] which is the first hierarchical protocol to perform clustering as it divides the nodes into number of clusters and in each cluster, one node is elected to be the cluster head that is responsible of sending to the BS. While the communication distance is reduced in LEACH compared to direct routing, the distance between each some cluster heads and the BS will still be long enough to highly consume the sensor's energy. Moreover, the cluster heads are randomly selected regarding their remaining energy, which will affect the robustness and the lifetime of the network. In addition, LEACH perform the clustering in each transmission routing which produce high overhead. Other
hierarchical known protocol is HEED Protocol [3]. It extends the basic scheme of LEACH by using the residual energy of the node as primary parameter for cluster head selection, and the network topology features such as node degree, distances to neighbors are used as secondary parameters to break the tie between the candidate cluster heads. However, same as LEACH protocol, HEED suffers from high overhead due the repeated clustering process in each round. Moreover, the nodes near the BS will die sooner due huge workload [4].

To overcome some of the main limitations in the previous protocols, new clustering approach were suggested in PEGASIS protocol [5].

1.1. PEGASIS protocol

The protocol proposed by Lindsey and Raghavendra in 2002 as the first hierarchical routing protocol to perform chain-based clustering. The goal of their work was to save the sensor nodes energy consumption regardless of the other consequences on the network performance. However, many research have been done to improve the chain-based routing and overcome the deficiencies in PEGASIS protocol such as in [6-11]. One of these improvements was by IEEPB protocol. Table 1 shows a summarization of the characteristics of some of the recent protocols.

| Ref. | Protocol          | Chain formation method | No. of chains | Leader node selection method | Proposed method                                                                 |
|------|-------------------|------------------------|---------------|------------------------------|---------------------------------------------------------------------------------|
| [6]  | Hadjila et al.    | the nodes are sorted based on their ordinates | Multiple      | The closest node to the chain is the leader node | Form multiple parallel chains in the direction of the sink. Form a main chain includes first nodes of each chain. |
| [7]  | Gupta and Saraswat | Greedy algorithm with the sensor nodes allowed to opt visited nodes again if they are the nearest one from those sensor | Single        | Considers the degree of nodes in addition to the residual energy and the distance to the BS | At the chain formation phase, the end node connects to its nearest node with the constraint that it can only connects to a node closer to the BS. |
| [8]  | Ruan et al.       | Ant Colony Optimization [13] | Multiple      | Considers the residual energy of the node and the distance to the BS | Uses the neural network to select the leader nodes. |
| [9]  | SCBC              | Greedy algorithm starting by formalizing two chains at once that later connected together. | Multiple      | Considers the residual energy and cost function. Secondary leader nodes are selected to send to the BS. | Divide the sensing are into sectors to form multiple chains. In order to reduce the delay accused by long chains. |
| [10] | Ghosh et al.      | Ant Colony Optimization [13] | Single        | Considers both the residual energy of the node and the distance to BS | Distributed Dominating Set Formation (DDSF) to choose the action nodes, the rest will be in a sleep mode. |
| [11] | Jawad and Ali     | Greedy algorithm       | Multiple      | Considers both the residual energy of the node and the distance to BS | Proposed using the k-means clustering algorithm to divide the nodes. |
| [12] | IEEPB             | Advanced greedy algorithm | Single        | Considers both the residual energy of the node and the distance to BS | Proposed an additional step to the greedy algorithm to reduce the distance between the connected nodes. |
| [13] | CHIRON            | Greedy algorithm       | Multiple      | Considers the residual energy only. | Proposed dividing the sensing area into leaves where the BS is the central point. |
| [14] | Shekh et al.      | Greedy algorithm       | Multiple      | Considers the nodes density in addition to the residual energy and distance from the BS | Proposed dividing the sensing moreover CHIRON protocol by dividing each leaf into equal sizes groups to have more similar in length chains. |
| [15] | Patel and Munjani | Greedy algorithm       | Multiple      | In turn. All the alive nodes have equal chance to be the leader node. Same as PEGASIS | Proposed multi-chain PEGASIS where the only difference from pegasus is the sensing area is divided into four regions |
| [16] | Bhatti and Rama   | Modified Greedy algorithm with fuzzy system | Single        | Based on the measured RSSI value | Modifies the PEGASIS protocol using fuzzy system and cuckoo search algorithm for optimization |
| [17] | PEGASIS-INL       | Greedy algorithm       | Single        | Based on the measured RSSI value | A subset of the nodes selected to be candidate leaders and multiple-overlapped chains formed with the candidate leaders as root. |
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1.2. IEEPB protocol

An Improved Energy-Efficient PEGASIS-Based Protocol (IEEPB) is proposed by Feng et al. in 2011 [12]. The protocol main contribution is modifying the chain forming algorithm to farther minimize the transmission energy consumption. Although the results in [12] showed an improved performance of IEEPB, we wanted to test the protocol performance under several scenarios in comparison with PEGASIS.

The next section will discuss the methods used in the chain formation phase for the PEGASIS protocol and for IEEPB protocol, followed by the leader node selection methods in section III. To compare the performance, Section IV presents the simulation environment and the results. Finally, a discussion and analysis of the results provided in Section V and the conclusion in Section VI.

2. RESEARCH METHOD

2.1. Chain Formation Methods

In this stage of the chain-based protocols, nodes are set to form a chain-like topology between them; therefore, when the data transmission phase comes each node will send its sensed and aggregated data to their neighbor node in the chain, in the direction of the leader node of the chain. Many algorithms proposed to perform the chain formation phase in the chain-based protocols and the process could be centralized, i.e. processed at the BS, or distributed, i.e. processed by the nodes themselves through communication among them.

2.1.1. PEGASIS protocol

In PEGASIS, they follow the greedy approach in the chain formation process where each node connects to its nearest node that did not join the chain yet. The chain formation starts from the farthest node from the BS, in order to make sure that nodes far from the BS have close neighbors, since the neighbor distance in the greedy approach will increase gradually as more nodes join the chain and less nodes left available to connect. Figure 1 shows a chain of 200 nodes formed using the greedy algorithm as in PEGASIS protocol. In case of a node dies, the whole chain is reformed again in the same manner.
2.1.2. IEEPB protocol

An enhanced greedy algorithm for the chain formation process was adopted in IEEPB, where the comparison of distance occurs twice before any node join the chain. The chain formed by the enhanced algorithm as follow:
1. The node farthest from the BS join the chain first and labeled as the end node of the chain
2. End node of the chain finds the nearest node that did not join the chain yet and sets it as next node waiting to join the chain.
3. Next node compare distance from itself to nodes already on the chain and connects with the nearest node of them.
4. After joining the chain, the next node becomes the end node of the chain and steps 2-4 repeated until all nodes join the chain.

Similar to PEGASIS, in case of a node dies the whole chain is reformed again in order to bypass the dead node. Figure 2 shows a chain of 200 nodes formed using the enhanced greedy algorithm as in IEEPB protocol.

2.2. Leader Selection Method

The leader node in HRP is the node responsible of forwarding the received data to the BS. Therefore, choosing the most fitting node to be a leader node is an important step:
2.2.1. PEGASIS protocol
Nodes in PEGASIS take turns to be the leader of the round. In round \(i\), the node number \(i \mod N\) is the leader, where \(N\) is total number of nodes. There is no consideration of how far the node located from the BS nor any other parameter; all nodes will have equal chance to become a leader node.

2.2.2. IEEPB protocol
For more suitable leader selection, the IEEPB uses weighting method that considers both the residual energy of the node and its distance from BS as parameters. In each round, the combined weight of each node is compared and the node with minimum weight, i.e. the less costly to be the leader of the round is selected as the leader node of the round.

3. RESULTS AND ANALYSIS
3.1. Simulation Environment
The simulation of both protocols done using MATLAB 2017a [24] where \(N\) sensor nodes randomly distributed on a square sensing field with the BS fixed far from the sensing field.

3.2. Energy Model
The same radio model described in [25] is adopted in this simulation. In this model, to transmit an \(L\) bit message for a distance \(d\), the radio expends:

\[
E_T(L, d) = E_{elec} \times L + E_{amp} \times L \times d^n
\]

(1)
to receive an \(L\) bit message, in the model the radio expends:

\[
E_R(L) = E_{elec} \times L
\]

(2)

Where \(E_{elec}\) is the energy dissipated per bit to run the transmitter or the receiver circuits and the \(E_{amp}\) is the energy dissipation of the transmission amplifier depending on the distance to the receiver. In equation 1, if the distance between transmitter node and receiver node is less than a threshold distance then the free space channel model is used where \((n=2, E_{amp}=E_{fs})\); otherwise multipath fading channel model is used where \((n=4, E_{amp}=E_{mp})\). Table 2 presents the network parameters used in simulation of both protocols.

| Parameter             | Value                  |
|-----------------------|------------------------|
| Number of nodes (\(N\)) | 200 and 400            |
| Network size          | 100 m * 100 m          |
| BS location           | (50, 175)              |
| Initial energy        | 0.5 J                  |
| \(E_{elec}\)          | 50 nJ/bit              |
| \(E_{fs}\)            | 100 pJ/bit/m^2         |
| \(E_{mp}\)            | 0.0013 pJ/bit/m^4      |
| The threshold distance| \((E_{fs}/E_{mp})\)     |
| Data packet size      | 2000 bit               |
| Energy of data aggregation EDA | 5 nJ/bit |

Table 2. Simulation parameters

3.3. Results of Simulation
The results are average of at least five simulation times. Figure 3 shows the network lifetime of 200 nodes by presenting the number of alive nodes from first round until last node dies. In PEGASIS, the average round the first node dies was round 548 while in IEEPB is round 1160 with 53% improvement. Moreover, the network dies at average 1633 rounds in PEGASIS and at 1831 rounds in IEEPB with 11% improvement. In addition, Figure 4 shows the total residual energy of all nodes throughout the network lifetime.

Same pervious performance metrics applied on 400 nodes network and results in as Figures 5 and 6 that shows PEGASIS protocol giving similar performance as in 200 network while the performance decreased clearly with IEEPB protocol.

In Figure 5, the 53% improvement in the average round the first node dies in IEEPB over PEGASIS decreased to 34% and the whole network dies with 6% improvement over PEGASIS instead of 11% as in with 200 nodes network. Moreover, Figure 6 shows IEEPB protocol residual energy freely dropping around the round 1400.
Another metric we considered is the overhead. We measured the total number of control packets exchanged between the sensor nodes and the BS to set up the routing path including chain formation and reformations throughout the network lifetime. Figure 7 shows the overhead for both 200 and 400 nodes networks.

3.4. Analysis of Results

In general, the greedy algorithm used in PEGASIS causes long link (LL) problem in the chain, as shown in Figure 1, because of the last nodes joining the chain will have fewer choices to connect to, since nodes in greedy approach cannot be revisited. However, in the enhanced greedy algorithm in IEEPB nodes
can be revisited, i.e. nodes are allowed to have more than two connections, as shown in Figure 2, which solved the LL problem in PEGASIS. Therefore, solving the LL problem means larger sending distance between nodes are avoided which improved the energy consumption of the nodes compared to the greedy algorithm.

Based on our experiments of different networks sizes simulation, the avoidance of LL problem showed obvious improvement in performance of networks with less nodes density, as in Figure 3 with 200 nodes; however, as the density of nodes increased the avoiding of LL problem was not enough for IEEPB to give better performance than PEGASIS. For instance, Figure 5 shows the lifetime of 400 nodes network in the same 100*100 m sensing area and IEEPB performance clearly decreased compared to Figure 3. In addition, at some point in Figure 5 when more than 80 nodes were alive, the total number of alive nodes in PEGASIS protocol was more than IEEPB protocol.

This fell in IEEPB performance is due the increase of the probability of a node to have more connected nodes (child node) as the number of nodes in area increase. Therefore, the energy consumption of each node will be larger in IEEPB since nodes consume $E_{th}(L)$ per every child node while in PEGASIS each node has only one child node to receive data from. This problem could be solved by forcing a constraint on the number of connected nodes, or by decreasing the density of nodes by having multiple chains in a network instead of single chain.

In this paper, we assumed a distributed chain formation process; therefore, nodes have to communicate by exchanging control packets in order to form the chain. In Figure 7, the results show that PEGASIS has total number of control packets less than IEEPB; this is due the greedy algorithm in PEGASIS requiring less communications for a node to join the chain in comparison of the enhanced greedy algorithm in IEEPB.

Moreover, the overhead caused by the enhanced greedy algorithm increase as the number of nodes involved in a chain formation process increase, therefore, reducing the overhead could achieved by having multiple $K$ chains as the communication will minimize $K$ times ($K$ is number of chains).

4. CONCLUSION

In conclusion, the enhanced greedy algorithm used by IEEPB protocol solved the LL problem caused by the greedy algorithm in PEGASIS, therefore, the energy consumption for data transmission is reduced. Moreover, the consideration of residual energy and distance from the BS in IEEPB improved the energy consumption as well. However, the simulation results showed a fallback in IEEPB performance including the network lifetime, energy consumption, and total overhead, as the nodes the number of deployed increases, i.e. higher density. Therefore, a discussion of possible solutions of the problem is presented. Our future work will focus on improving the IEEPB protocol and solve the limitations in case of a high-density network.

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