**μ-Planner: A Robot Path Planning Approach Based on Language Measure of Unsupervised Automata**

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**Abstract:** This paper proposes a robot path planner based on language measure, \( \mu \)-planner. Workspace is discretized in an occupancy grid map and we model the system by considering how events, associated to robot’s motions, take it to different cells (discrete positions). The calculated language measure values corresponds to a gradient, which the robot can use reach its destination by choosing events that take it to states with higher measure values. Concepts of Laplace’s equation and harmonic functions are used to prove that our method guarantees both the existence and monotonicity of language measure. The proposed method is simple and computationally inexpensive and guarantees existence of path from any co-accessible state to the destination. Experiments considering different scenarios have been performed to validate and compare \( \mu \)-planner with similar methods.

**Keywords:** Path planning, Language measure, Event probability, Mobile robots and grid map.

1. **INTRODUCTION**

Path planning is one of the most basic tasks to be performed in mobile robot applications. Several methods have been proposed in the last decades to allow robots calculating a path from its current position to a desired destination [10-13, 16, 19].

Based on the mathematical tools used to describe the workspace and calculate a path taking the robot toward the destination, most methods can be classified in grid maps, roadmaps and potential field planners, [17]. Path planners based on both grid maps and roadmaps describe how a robot can move through the workspace using a graph, than search algorithms, such as Dijkstra and \( A^* \), can be used to obtain a path connecting the robot’s initial position to the destiny. The main difference in these methods is how they discretize workspace in nodes and define edges connecting them.

Potential field methods define a potential value for each robot configuration (position in the workspace, for instance). Potential values usually are calculated defining repulsive forces from obstacles and attractive forces from destiny.

Authors also proposed path planning methods based on Discrete Event Systems (DES) [2, 18, 24]. Most methods use automaton or Petri net structures and formal verification to generate paths, as sequences of events representing robot actions or states representing discrete positions in the workspace. Recently, methods based on language measure, [4, 5, 21, 22], have been proposed to calculate robot paths. The basic idea of language measure is to attribute a value to each state based on how close they are to marked states and how many event strings intersects on them.

Path planners based on language measure model the robot possible motions through a workspace as an automaton, marking states representing the destiny and obstacles. By attributing positive values (+1) to the destiny state and negative (-1) to the obstacles (or collision) states, the methods can generate a gradient without local maximum (or minimum), allowing the robot to reach the destination from any co-accessible state. However, the computational cost can be quite high.

This paper proposes \( \mu \)-planner, a simple and computationally efficient method based on language measure. The method is able to produce paths similar to those obtained from [4 and 5] at much lower computational cost. \( \mu \)-planner does not require a supervisory control to ensure the existence of numerical solution and absence of local minimum or maximum values. By defining an automaton structure that already guarantees both existence of a solution and its monotonicity, \( \mu \)-planner avoids iterative processes to obtain a supervisory, being able to calculate language measure with a single matrix inversion.

2. **LANGUAGE MEASURE THEORY**

Let \( G = (Q, \Sigma, \delta, q_{init}, Q_m) \) be a deterministic finite-state automaton (DFSA) and \( \hat{\delta} : Q \times \Sigma^* \rightarrow Q \) the extended state transition function. In addition, let \( G_i \) be a version of \( G \) in which \( q_{init} = q_i \in Q \).
Definition 1. The generated and marked languages of an automaton $G_i$, $L(G_i)$ and $L_m(G_i)$, are defined as:

$$L(G_i) = \{ s \in \Sigma \mid \delta(q_i,s) \in Q \}$$

(1)

$$L_m(G_i) = \{ s \in \Sigma' \mid \delta(q_i,s) \in Q_m \}$$

(2)

Definition 2. Language $L(q_i,q_j)$ (or simply $L_{ij}$) corresponds to the set of strings starting in $q_i$ and terminating in $q_j$. Formally, $L(q_i,q_j)$ is defined as:

$$L(q_i,q_j) = \{ s \in \Sigma' \mid \delta(q_i,s) = q_j \in Q \}$$

(3)

Language $L(q_j)$, the set of all strings allowed in $G$ from state $q_j$, is defined as:

$$L(q_j) = \bigcup_{q_i \in Q} L(q_i,q_j) = L(G_i)$$

(4)

Language $L(G)$ corresponds to the set with all strings initiating at any state of $Q$. Formally, we have that:

$$L(G) = \bigcup_{q_i \in Q} L(q_i)$$

(5)

The set of marked states, $Q_m$, can be partitioned as $Q_m = Q_m^c \cup Q_m^\prime$, where $Q_m^c$ represents the set of states we desire to reach and $Q_m^\prime$ the states we have to avoid.

Definition 3. The language measure function $\mu : 2^{L(G)} \to \mathbb{R}$ associates a real value to a language $L(q_j) \in L(G)$, such that:

$$\mu(L(q_j)) = \begin{cases} 
0, & q_j \notin Q_m^c \\
> 0, & q_j \in Q_m^c \\
< 0, & q_j \in Q_m^\prime 
\end{cases}$$

(6)

In order to give a physical meaning to the measure $\mu$ of a language $L(q_j)$, [22] defined it as the probability of reach a state in $Q_m^c$ from a state $q_j$, while avoiding states in $Q_m^\prime$.

The computation of $\mu$ relies on three structures: event probability (or cost) matrix $\tilde{\Pi}$; state transition probability matrix $\Pi$ and the characteristic function $\chi : Q \to [-1,1].$

Definition 4. For each event $\sigma_k \in \Sigma$ and state $q_j \in Q$, the probability of triggering $\sigma_k$ at $q_j$, $\tilde{\pi}(q_j,\sigma_k)$ (or simply $\tilde{\pi}_k$), is defined such that:

1. $\tilde{\pi}_k \in [0,1)$ and $\sum_k \tilde{\pi}_k < 1$
2. $\tilde{\pi}(q_j,\epsilon) = 1$ and $\tilde{\pi}_k = 0$ if $\delta(q_j,\sigma_k)$ is undefined;
3. $\tilde{\pi}(q_j,\sigma_k,s) = \tilde{\pi}(q_j,\sigma_k)\tilde{\pi}(\delta(q_j,\sigma_k),s)$.

Item (1) provides a sufficient condition to existence of a finite value of $\mu$ [22]. Items (2) and (3) provide an iterative way to get the probability of occurring a string $s$ from a state $q_j$.

Considering $|Q| = n$ and $|\Sigma| = l$, the event probability matrix is defined as:

$$\tilde{\Pi} = \begin{bmatrix} 
\tilde{\pi}_{11} & \tilde{\pi}_{12} & \cdots & \tilde{\pi}_{1l} \\
\tilde{\pi}_{21} & \tilde{\pi}_{22} & \cdots & \tilde{\pi}_{2l} \\
\vdots & \vdots & \ddots & \vdots \\
\tilde{\pi}_{nl} & \tilde{\pi}_{n2} & \cdots & \tilde{\pi}_{nl} 
\end{bmatrix}$$

(7)

Definition 5. The probability of reaching a state $q_j$ from $q_i$, with the occurrence of a single event, is defined as:

$$\pi_{ij} = \sum_{\sigma \in \Sigma | \delta(q_i,\sigma) = q_j} \tilde{\pi}(q_i,\sigma)$$

(8)

Based on (8), the state transition probability matrix is defined as $\Pi_{ij} = \pi_{ij}$.

Definition 6. The characteristic function $\chi : Q \to [-1,1]$ allows the designer to set weights on the states based on its perception of the application. Formally, we have that:

$$\forall q_i \in Q, \quad \chi(q_i) = \begin{cases} 
[-1,0), & q_i \in Q_m^- \\
(0,1], & q_i \in Q_m^+ \\
0, & q_j \notin Q_m 
\end{cases}$$

(9)

The state weighting vector, $\chi$ -vector, is defined by $\chi = [\chi_1, \chi_2, \ldots, \chi_n]^T$, where $\chi_i = \chi(q_i)$. The signed real measure of a language $L(q_j)$ is defined as:

$$\mu_i = \mu(L(q_j)) = \sum_{q_j \in Q} \mu(L(q_i,q_j))$$

(10)

where $\mu(L(q_i,q_j))$ is defined by (11).
\[ \mu(L(q_i,q_j)) = \sum_{q_i \in L(q_i,q_j)} \tilde{s}(q_i,s) \chi_j \] (11)

In [22], the authors show that Equation (10) can be expressed as:

\[ \mu_i = \sum_{q_j \in Q} \mu_j \cdot \chi_i \] (12)

Considering a matrix structure, Equation (12) can be expressed as:

\[ \mu = \Pi \mu + \chi \] (13)

The solution of (13) is given by:

\[ \mu = (I - \Pi)^{-1} \chi \] (14)

where \( I \) is the \( n \times n \) identity matrix.

After obtaining the \( \mu \)-vector, \( \mu = [\mu_1, \mu_2, \ldots, \mu_n] \), the system gets a metric from which it can choose the next action (enabled event in the plant \( G \)). By choosing the event that leads to the state with higher \( \mu \) value, the system will be executing the string with higher probability to reach a state in \( Q_m^+ \).

3. LANGUAGE MEASURE ON ROBOT PATH PLANNING

In the last decades, several works on path planning based on language measure have been proposed. Next, we present a brief overview of how language measure have been applied in robotics.

In [21 and 22], the authors define the basis of signed real language measure. Other works from the same research group address specific issues such as computational costs of the algorithms proposed to obtain the language measure, [20]; present proofs that, under some conditions, its always possible to obtain a finite measure \( \mu \), [23]; etc.

An important aspect often addressed by the authors is how guarantee that \( (I - \Pi) \) is an invertible matrix. In [23], the authors redefine \( \Pi \) based on Markov conditional probability and presents a method for estimating it that results in \( \Pi \) as a stochastic matrix.

As result of the use of stochastic matrices, the premise \( \sum_k \tilde{s}_{ik} < 1 \) is not satisfied anymore and there is no guarantee matrix \( (I - \Pi) \) is invertible. To circumvent such problem, the authors proposed choosing a convenient value of a parameter \( \theta \), such that \( 0 < \theta < 1 \), and calculating the language measure as:

\[ \mu(\theta) = (I - (1 - \theta)\Pi)^{-1} \chi \] (15)

The chosen \( \theta \) must be small enough to guarantee that \( \mu(\theta) \) is invariant to \( \theta \), i.e. \( \forall q_i,q_j \in Q \mid \mu_i < \mu_j \rightarrow \mu(\theta)_i < \mu(\theta)_j \), and higher enough to guarantee that there will be no numerical problem to calculate the inverse.

Chattopadhyay and Ray [6] addresses this problem and proposes an extension of language measure, the renormalized language measure. In [7], the authors proposed an algorithm to estimate the critical lower bound of \( \theta \), namely \( \theta^* \).

In [4 and 5], authors proposed a path planning method based language measure and supervisory control theories. The workspace is discretized in a grid and it’s considered the robot can move to one of its 8 neighbors, each motion represented by an event \( \sigma_k \in \Sigma_c \), with \( k = \{1, \ldots, 8\} \). Transitions to states representing occupied positions are allowed. However, once in such a state, there is a single uncontrollable event \( \sigma_u \) that can occur and taking plant \( G \) to \( O_b \) state, which represents a collision with an obstacle. Also, the boundaries of the workspace are not considered in the model, i.e. there are no states representing them.

Figure 1 illustrates how the workspace is discretized and an automaton \( G \) representing how the robot can move through the workspace grid cells. The events representing the robot motions where suppressed to simplify the figure.
Formally, the system is represented by an automaton \( G = (Q, \Sigma, \delta, \Gamma, q_{\text{init}}, Q_m) \). The set \( Q \) is formed by states representing both the free (\( Q_f \)) and occupied (\( Q_o \)) cells in the grid map and state \( O_b \), \( Q = Q_f \cup Q_o \cup \{O_b\} \). The alphabet is defined by \( \Sigma = \Sigma_c \cup \{\sigma_s\} \) and \( Q_m = \{q_{\text{goal}}, O_b\} \). Finally, function \( \Gamma: Q \rightarrow 2^\Sigma \) indicates which events are enabled at each state.

The event probability matrix, \( \prod \), is defined based on the number of events enabled in each state. For states representing occupied positions, only event \( \sigma_s \) is enabled and the probability is set as 1. Equation (16) presents how event probabilities are defined:

\[
\prod_{ij} = \begin{cases} 
1, & \text{if } \sigma_i \in \Gamma(q_i) \cap \Sigma_c \\
0, & \text{if } \sigma_i \notin \Gamma(q_i) \cap \Sigma_c \\
1, & \text{if } \sigma_i = \sigma_s \in \Gamma(q_i)
\end{cases}
\]  

(16)

Characteristic function \( \chi \) is defined as follows:

\[
\chi(q_i) = \begin{cases} 
-1, & \text{if } q_i \in Q_o \\
1, & \text{if } q_i = q_{\text{goal}} \\
0, & \text{otherwise}
\end{cases}
\]  

(17)

In order to ensure global monotonicity of the language measure, [4 and 5] propose an algorithm to compute the optimal supervisor for \( G \). Iteratively, the algorithm recalculates the language measure, defined by Equation (18), and use it to obtain the set of disabled transitions.

\[
\nu = \theta'[1 - (1 - \theta')\prod^{-1}\chi]
\]  

(18)

where \( \theta' \) is the critical lower bound of \( \theta \) and \( \prod \) is obtained according to Definition 5.

At each iteration, the algorithm disables controllable transitions \( q_i \rightarrow q_j \) such that \( \nu_i > \nu_j \). When a transition is disabled, the value \( \pi_y \) is added to \( \pi_x \) (self-loop) and \( \pi_y \) is set to 0. Then, \( \theta' \) and \( \nu \) are recalculated based on the updated \( \prod \). The algorithm terminates when the sets of disabled transitions \( \mathcal{B} \) of two consecutive iterations coincide. Final language measure, \( \nu_x \), is defined as the \( \nu \) from the last iteration.

Since the method ensures global monotonicity, the robot can navigate toward the goal by simply moving from its current position to its neighbor with the highest \( \nu \) value.

Other works focusing in specific aspects such as localization uncertainties, smoother paths, efficient replanning and navigation without global positioning facilities have been proposed, [3, 8, 9, 14 and 15]. However, as these works are based on the same framework to calculate \( \nu \), we do not present them in this paper.

4. PROPOSED METHOD

In this paper, the path planning problem is also addressed considering a language measure approach. However, we define event probability matrix \( \prod \) in a way that guarantees the existence of matrix \( (I - \prod) \) inverse and that \( \mu \), defined by Equation (14), corresponds to a globally monotonic gradient. Specifically, we propose a definition for event probability matrix \( \prod \) that results in a \( \mu \) function that holds maximum and minimum principle. By doing so, our method does not need to calculate neither a \( \theta \) value nor a supervisor, making it computationally inexpensive.

Maximum and minimum principle states that there is no local (or even global) maximum or minimum in the gradient inner region [1]. The maximum and minimum values occur only in the boundary region and critical points. In path planning applications, the boundary region, \( \partial \Omega \), corresponds to the workspace borders and obstacles. A single critical point represents the destination.

Let \( G = (Q, \Sigma, \delta, \Gamma, q_{\text{init}}, Q_m) \) be an automaton representing how the robot can move through a 2-D workspace and \( \Sigma = \{\sigma_1, ... , \sigma_s\} \) the event set, as illustrated in Figure 2. Also, let the set of states be defined as a partition \( Q = Q_f \cup Q_o \cup \{q_{\text{goal}}\} \), in which \( Q_o \) is the set of states representing positions in \( \partial \Omega \) and \( Q_f \) the states in the free space region. In this paper, we propose a \( \mu \) function definition, presented in Equation (19), that satisfies the maximum and minimum principle.

\[
\mu(q_i) = \begin{cases} 
\frac{1}{8} \sum_{q_i \in \mathcal{B}} \mu(\delta(q_i, \sigma_i)) , & \text{if } q_i \in Q_f \\
K_1 , & \text{if } q_i \in Q_o \\
K_2 , & \text{if } q_i = q_{\text{goal}}
\end{cases}
\]  

(19)
where \( K_1 \) and \( K_2 \) are chosen constants that just need to be different. If \( K_2 > K_1 \), the language measure grows toward the destiny.

**Proof.** The classical example of functions that holds maximum and minimum principle are the harmonic functions. Thus, the \( \mu \) defined by Equation (19) being a harmonic function is a sufficient, but not necessary, condition to ensure the maximum and minimum principle.

Let \( \ell: Q \rightarrow \mathbb{N}^2 \) be a function that maps a state \( q_i \) to its position \((x,y)\) in the grid cell and \( \mu(q_i) = \mu(x,y) \), if \( \ell(q_i) = (x,y) \).

A harmonic function is a solution for Laplace's equation, defined by Equation (20) for a 2-dimensional space:

\[
\nabla^2 U(q) = \frac{\partial^2 U}{\partial x^2} + \frac{\partial^2 U}{\partial y^2} = 0, \quad \forall q = (x,y) \in \Omega
\]

(20)

where \( U(q) \) is a harmonic function defined over region \( \Omega \).

To solve Laplace's equation in discretized environments, the partial differential equations are often replaced by finite difference equations, an approximation obtained by Taylor series around a point \((x_0,y_0)\).

By considering \( \mu(x,y) \) as a Taylor series limited to second-order (around \( x_0 \) and \( y_0 \)), we have that:

\[
\mu_{i,j} = \mu_{x_0,y_0} + (x-x_0) \frac{\partial \mu}{\partial x}(x_0,y_0) + (y-y_0) \frac{\partial \mu}{\partial y}(x_0,y_0) \\
+ \frac{(x-x_0)^2}{2} \frac{\partial^2 \mu}{\partial x^2}(x_0,y_0) + \frac{(y-y_0)^2}{2} \frac{\partial^2 \mu}{\partial y^2}(x_0,y_0) \\
+ (x-x_0)(y-y_0) \frac{\partial^2 \mu}{\partial x \partial y}(x_0,y_0)
\]

(21)

Evaluating (21) for \( \Delta_x > 0 \) and \( \Delta_y > 0 \) on both sides of \( x_0 \) and \( y_0 \), we can rewrite Equation (20) as a finite difference equation centered in \((x_0,y_0)\):

\[
\nabla^2 \mu(x,y) = \sum_{i,j \in [-1,1]} \sum_{|i|,|j| \\not= 0} \mu(x_0 + i\Delta_x, y_0 + j\Delta_y) - 8\mu(x_0,y_0) = 0
\]

(22)

Considering \( \Delta_x = \Delta_y = 1 \) (smallest displacement in the grid cell), Equation (22) can be rewritten as:

\[
\mu_{x+1,y-1} + \mu_{x+1,y} + \mu_{x+1,y+1} + \mu_{x,y} + \mu_{x-1,y} - 8\mu_{x,y} = 0
\]

(23)

Thus, for each position that is not a critical point or inside the boundary region, the \( \mu \) value can be calculated based on (23). Analysing Equation (23), one can notice it corresponds to the first line of our \( \mu \) definition, presented in Equation (19).

Considering Dirichlet's condition [1], we can set constant values to the potential of points in \( \partial \Omega \) and of critical points, such as:

\[
U(q) = \begin{cases} 
K_1, & q \in \partial \Omega \\
K_2, & q = q_{goal}
\end{cases}
\]

(24)

in which \( K_1, K_2 \in \mathbb{R} \) are chosen constant values.

Notice that Equation (24) corresponds to the remaining lines of Equation (19). Harmonic functions properties also guarantees that the magnitude of \(|K_2 - K_1|\) does not influence the gradient directions. Thus, for any \( K_2 > K_1 \), the generated path will be the same.

Our goal, then, is to propose a formulation to \( \vec{\Pi} \) such that state transition matrix \( \Pi \) and characteristic function \( \chi \) results in the \( \mu(q) \) definition from (19). To do so, we define automaton \( G \) transitions allowing only events taking the robot to a neighbor in \( Q_f \). Formally:

\[
\delta(q_i, \sigma_k) = \begin{cases} 
q_{j}, & \text{if } q_i, q_j \in Q_f \\
\text{undefined}, & \text{otherwise}
\end{cases}
\]

(25)

Based on automaton \( G \), \( \vec{\Pi} \) can be defined as:

\[
\vec{\Pi}_{i,k} = \begin{cases} 
\frac{1}{8}, & \text{if } \sigma_k \in \Gamma(q_i)0.1cm \\
0, & \text{otherwise}
\end{cases}
\]

(26)

The characteristic function \( \chi \) is defined as follows:
\( \chi(q_i) = \begin{cases} K_1, & \text{if } q_i \in Q_o \\ K_2, & \text{if } q_i = q_{goal} \\ 0, & \text{otherwise} \end{cases} \) (27)

Thus, the system defined by Equation (19), for all \( q \in Q \), can be represented by:

\[ \mu = \mu \prod + \chi \] (28)

As presented in section 2, \( \mu \) can be obtained as:

\[ \mu = (1 - \prod)^{-1} \chi \] (29)

However, our automaton \( G \) and \( \prod \) definitions guarantees both the existence of a numerical solution for the system and monotonicity of the obtained \( \mu \).

The path taking the robot from its current position to destination can be obtained as sequence of neighbor states with highest \( \mu \) value. Formally:

\[ \text{path} = [q_1, q_2, \ldots, q_n] \] (30)

where \( q_1 = q_{init} \), \( q_n = q_{goal} \), and \( q_{j+1} = \delta(q_j, \sigma) \), such that \( \sigma = \max_{\alpha \in \mathcal{C}(q_j)} \mu(\delta(q_j, \sigma_j)) \).

Next, we present the experiments performed in order to evaluate the similarity of the plans obtained from \( \mu \)-planner with those obtained using \( \nu \) path planning.

5. RESULTS AND DISCUSSIONS

In order to validate the proposed method and evaluate the quality of paths it generates, we perform experiments considering different workspaces and robot positions. The path planner proposed by [5], \( \nu \)’, is also implemented and used to generate paths in the same scenarios. Both methods are compared based on the time necessary to calculate the language measure vectors and the similarity of the paths. All experiments were performed using a PC with 4 GB of RAM and Processor Intel I3 running Linux Mint (18.2). Next, we present the metrics used to evaluate the paths.

5.1. Metrics

Paths are compared regarding four metrics: number of steps, length and minimum and average distance to obstacles. Number of steps, \( n_{steps} \), corresponds to the number states in the path, while the length, \( P_{dis} \), is the real distance the robot have to move. These metrics can be defined as:

\[ n_{steps} = |\text{path}| - 1 \] (31)

\[ P_{dis} = \sum_{i=1}^{n_{steps}} \|\ell(q_i) - \ell(q_{i-1})\| \] (32)

where \( q_i \in \text{path} \) is the \( i \)-th state in the path, \( \ell(q_i) \) corresponds to \( q_i \) position in the grid and operator \( \| \| \) is the Euclidean distance.

Regarding the distance to obstacles, let \( d_{co}(q_i) \) be the Euclidean distance between \( q_i \)'s cell (position in the grid map) and the closest obstacle. The minimum, \( d_{min} \), and average, \( \bar{d} \), distances to obstacles can be defined as:

\[ d_{min} = \min_{q_i \in \text{path}} d_{co}(q_i) \] (33)

\[ \bar{d} = \frac{1}{|\text{path}|} \sum_{q_i \in \text{path}} d_{co}(q_i) \] (34)

Figure 3 illustrates part of a path and the distance between each cell in the path and its closest obstacle.

5.2. Experiments

We performed experiments in three different workspaces, presented in Figure 4. For each workspace, a destiny and 10 different initial positions were randomly chosen. Then, the proposed method and \( \nu \)’ path planner were used to calculate paths for each configuration. For simplicity, we chose \( K_1 = -1 \) and \( K_2 = 1 \) for our method.

Additionally, we calculate paths using \( \nu \)’ algorithm with constant (arbitrary) values of \( \theta \) (10\(^{-2}\) and 10\(^{-3}\)). Thus, one can observe the impact of \( \theta \) values in path generation and the cost of calculating \( \theta^* \) (lower bound of \( \theta \)) at each iteration.
Figures 5, 6 and 7 present the paths, for a single \((q_{\text{init}}, q_{\text{goal}})\) configuration, generated by the proposed method and \(\nu^*\).

Table 1: Results of Experiments in a Workspace Prone to Local Minimum

|               | \(t_{\text{exp}}\) (s) | \(n_{\text{exp}}\) | \(p_{\text{ass}}\) | \(d_{\text{ass}}\) | \(\bar{d}\) |
|---------------|----------------------|-----------------|----------------|----------------|---------|
| Proposed     | 0.084 ± 0.001        | 21              | 28.04          | 1              | 1.34 ± 0.50 |
| \(\nu^*\)    | 1.590 ± 0.025        | 23              | 30.87          | 1              | 1.27 ± 0.47 |
| \(\nu^* (\theta = 10^{-3})\) | 0.995 ± 0.098        | 23              | 30.87          | 1              | 1.27 ± 0.47 |
| \(\nu^* (\theta = 10^{-2})\) | 0.957 ± 0.031        | 21              | 28.87          | 1              | 1.25 ± 0.46 |

5.3. Discussions

Both Figures 5, 6 and 7 and Tables 1, 2, 3 and 4 show the proposed method generate paths similar to those obtained using \(\nu^*\) path planner. For instance, in experiments with the workspace prone to local minimum, the proposed method generates paths similar to those obtained using \(\nu^*\) path planner.
minimum, the difference in the number of steps is $1 \pm 1$. Such value represents less than 0.5% of $n_{\text{steps}}$, for the paths presented in Figure 5. Regarding the distance to obstacles, the difference is less than 1 cell unit. Differences increases a little for workspaces with scattered obstacles, due to the higher number of possible paths, but it still less than 1% of the values presented in Table 3.

On the other hand, Tables 1, 2 and 3 show that, as the workspace’s size increases, $\nu'$ path planner becomes computationally expensive due to $\theta'$ calculation. Notice that, when $\theta'$ is previously defined, the time necessary to execute $\nu'$ decreases significantly.

Another important aspect is that, for both the proposed method and $\nu'$ path planner, diagonal motions “cost” the same as an one direction motion (up, down, left or right). So, “unnecessary” diagonal motions may occur in the generated paths, as in $\nu'$’s path shown in Figure 7. However, by defining $\Pi$ matrix based on Laplace’s equation, the proposed method generates smoother paths (as can be better illustrated in Figure 6), usually keeping farther away from obstacles.

### 6. CONCLUSIONS

This paper presents a path planning method based on language measure of unsupervised automata. By considering path planning problem from both DES and the maximum and minimum principle viewpoints, we propose a methodology to define the event probability matrix, $\Pi$, that guarantees existence of a numerical solution for $\mu = (I - \Pi)^{-1} \chi$.

Additionally, $\Pi$’s definition guarantees the monotonicity of the solution. Thus, language measure $\mu$ can be view as a gradient without local maximum (or minimum) and, from any free cell, it’s possible to generate a path to destiny by iteratively choosing the neighbor with highest $\mu$ value. Also, changes in the workspace can be handled online by adding/removing transitions in automaton $G$ and recalculating $\mu$, since $\mu$-planner is computationally inexpensive.
Experiments with several workspaces show $\mu$-planner generates paths similar to those obtained using $\nu$-planner. However, since our definitions already guarantee existence of solution, there is no need to iteratively compute a supervisor and $\theta$ values, which decreases greatly our method computational cost. Specifically, $\mu$-planner considers a single matrix inversion to calculate $\mu$ language measure vectors.

6.1 Future Works

In future works, strategies to smooth generated paths and better cope with dynamic environments, as those proposed in [3 and 15], will be addressed. Additionally, other definitions will be studied, in order to allow considering different “costs” for diagonal motions.

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