Interpolation of inverse operators for preconditioning parameter-dependent equations *

Olivier ZAHM† and Anthony NOUY†‡

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Abstract

We propose a method for the construction of preconditioners of parameter-dependent matrices for the solution of large systems of parameter-dependent equations. The proposed method is an interpolation of the matrix inverse based on a projection of the identity matrix with respect to the Frobenius norm. Approximations of the Frobenius norm using random matrices are introduced in order to handle large matrices. The resulting statistical estimators of the Frobenius norm yield quasi-optimal projections that are controlled with high probability. Strategies for the adaptive selection of interpolation points are then proposed for different objectives in the context of projection-based model order reduction methods: the improvement of residual-based error estimators, the improvement of the projection on a given reduced approximation space, or the recycling of computations for sampling based model order reduction methods.

1 Introduction

This paper is concerned with the solution of large systems of parameter-dependent equations of the form

\[ A(\xi)u(\xi) = b(\xi), \]

where \( \xi \) takes values in some parameter set \( \Xi \). Such problems occur in several contexts such as parametric analyses, optimization, control or uncertainty quantification, where \( \xi \) are random variables that parametrize model or data uncertainties. The efficient solution of equation (1) generally requires the construction of preconditioners for the operator \( A(\xi) \),

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either for improving the performance of iterative solvers or for improving the quality of residual-based projection methods.

A basic preconditioner can be defined as the inverse (or any preconditioner) of the matrix $A(\xi)$ at some nominal value parameter value $\bar{\xi} \in \Xi$ or as the inverse (or any preconditioner) of a mean value of $A(\xi)$ over $\xi$ (see e.g. [17, 16]). When the operator only slightly varies over the parameter set $\Xi$, these parameter-independent preconditioners behave relatively well. However, for large variabilities, they are not able to provide a good preconditioning over the all parameter set $\Xi$. A first attempt to construct a parameter-dependent preconditioner can be found in [14], where the authors compute though quadrature a polynomial expansion of the parameter-dependent factors of a LU factorization of $A(\xi)$. More recently, a linear Lagrangian interpolation of the matrix inverse has been proposed in [10]. The generalization to any standard multivariate interpolation method is straightforward. However, standard approximation or interpolation methods require the evaluation of matrix inverses (or factorizations) for many instances of $\xi$ on a prescribed structured grid (quadrature or interpolation), that becomes prohibitive for large matrices and high dimensional parametric problems.

In this paper, we propose an interpolation method for the inverse of matrix $A(\xi)$. The interpolation is obtained by a projection of the inverse matrix on a linear span of samples of $A(\xi)^{-1}$ and takes the form

$$P_m(\xi) = \sum_{i=1}^{m} \lambda_i(\xi) A(\xi_i)^{-1},$$

where $\xi_1, \ldots, \xi_m$ are $m$ arbitrary interpolation points in $\Xi$. A natural interpolation could be obtained by minimizing the condition number of $P_m(\xi) A(\xi)$ over the $\lambda_i(\xi)$, which is a Clarke regular strongly pseudoconvex optimization problem [32]. However, the solution of this non standard optimization problem for many instances of $\xi$ is intractable and proposing an efficient solution method in a multi-query context remains a challenging issue. Here, the projection is defined as the minimizer of the Frobenius norm of $I - P_m(\xi) A(\xi)$, which is a quadratic optimization problem. Approximations of the Frobenius norm using random matrices are introduced in order to handle large matrices. These statistical estimations of the Frobenius norm allow to obtain quasi-optimal projections that are controlled with high probability. Since we are interested in large matrices, $A(\xi_i)^{-1}$ are here considered as implicit matrices for which only efficient matrix-vector multiplications are available. Typically, a factorization (e.g. LU) of $A(\xi)$ is computed and stored. Note that when the storage of factorizations of several samples of the operator is unaffordable or when efficient preconditioners are readily available, one could similarly consider projections of the inverse operator on a linear span of preconditioners of samples of the operator. However, the resulting parameter-dependent preconditioner would be no more an interpolation of preconditioners. This straightforward
extension of the proposed method is not analyzed in the present paper.

The paper then presents several contributions in the context of projection-based model order reduction methods (e.g. Reduced Basis, Proper Orthogonal Decomposition (POD), Proper Generalized Decomposition) that rely on the projection of the solution $u(\xi)$ of (1) on a low-dimensional approximation space. We first show how the proposed preconditioner can be used to define a Galerkin projection based on the preconditioned residual, which can be interpreted as a Petrov-Galerkin projection of the solution with a parameter-dependent test space. Then, we propose adaptive construction of the preconditioner, based on an adaptive selection of interpolation points, for different objectives: (i) the improvement of error estimators based on preconditioned residuals, (ii) the improvement of the quality of projections on a given low-dimensional approximation space, or (iii) the recycling of computations for sample-based model order reduction methods. Starting form a $m$-point interpolation, these adaptive strategies consist in choosing a new interpolation point based on different criteria. In (i), the new point is selected for minimizing the distance between the identity and the preconditioned operator. In (ii), it is selected for improving the quasi-optimality constant of Petrov-Galerkin projections which measures how far the projection is from the best approximation on the reduced approximation space. In (iii), the new interpolation point is selected as a new sample determined for the approximation of the solution and not of the operator. The interest of the latter approach is that when direct solvers are used to solve equation (1) at some sample points, the corresponding factorizations of the matrix can be stored and the preconditioner can be computed with a negligible additional cost.

The paper is organized as follows. In section 2 we present the interpolation method for the inverse of a parameter dependent matrix. In Section 3 we show how the preconditioner can be used the definition of a Petrov-Galerkin projection of the solution of (1) on a given reduced approximation space, and we provide an analysis of the quasi-optimality constant of this projection. Then, different strategies for the selection of interpolation points for the preconditioner are proposed in Section 4. Finally, in Section 5 numerical experiments will illustrate the efficiency of the proposed preconditioning strategies for different projection-based model order reduction methods.

Note that the proposed preconditioner could be also used (a) for improving the quality of Galerkin projection methods where a projection of the solution $u(\xi)$ is searched on a subspace of functions of the parameters (e.g. polynomial or piecewise polynomial spaces) [13, 26, 28], or (b) for preconditioning iterative solvers for (1), in particular solvers based on low-rank truncations that require a low-rank structure of the preconditioner [24, 27, 18, 19]. These two potential applications are not considered here.
2 Interpolation of the inverse of a parameter-dependent matrix using Frobenius norm projection

In this section, we propose a construction of an interpolation of the matrix-valued function \( \xi \mapsto A(\xi)^{-1} \in \mathbb{R}^{n \times n} \) for given interpolation points \( \xi_1, \ldots, \xi_m \) in \( \Xi \). We let \( P_i = A(\xi_i)^{-1} \), \( 1 \leq i \leq m \). For large matrices, the explicit computation of \( P_i \) is usually not affordable. Therefore, \( P_i \) is here considered as an implicit matrix and we assume that the product of \( P_i \) with a vector can be computed efficiently. In practice, factorizations of matrices \( A(\xi_i) \) are stored.

2.1 Projection using Frobenius norm

We introduce the subspace \( Y_m = \text{span}\{P_1, \ldots, P_m\} \) of \( \mathbb{R}^{n \times n} \). An approximation \( P_m(\xi) \) of \( A(\xi)^{-1} \) in \( Y_m \) is then defined by

\[
P_m(\xi) = \arg\min_{P \in Y_m} \|I - PA(\xi)\|_F,
\]

where \( I \) denotes the identity matrix of size \( n \), and \( \| \cdot \|_F \) is the Frobenius norm such that \( \|B\|_F^2 = \langle B, B \rangle_F = \text{trace}(B^T C) \). Since \( A(\xi_i)^{-1} \in Y_m \), we have the interpolation property \( P_m(\xi_i) = A(\xi_i)^{-1} \), \( 1 \leq i \leq m \). The minimization of \( \|I - PA\|_F \) has been first proposed in [21] for the construction of a preconditioner \( P \) in a subspace of matrices with given sparsity pattern (SPAI method). The following proposition gives some properties of the operator \( P_m(\xi)A(\xi) \) (see Lemma 2.6 and Theorem 3.2 in [20]).

**Proposition 2.1** Let \( P_m(\xi) \) be defined by (2). We have

\[
(1 - \alpha_m(\xi))^2 \leq \|I - P_m(\xi)A(\xi)\|_F^2 \leq n(1 - \alpha_m^2(\xi)),
\]

where \( \alpha_m(\xi) \) is the lowest singular values of \( P_m(\xi)A(\xi) \) verifying \( 0 \leq \alpha_m(\xi) \leq 1 \), with \( P_m(\xi)A(\xi) = I \) if and only if \( \alpha_m(\xi) = 1 \). Also, the following bound holds for the condition number of \( P_m(\xi)A(\xi) \):

\[
\kappa(P_m(\xi)A(\xi)) \leq \frac{\sqrt{n - (n - 1)\alpha_m^2(\xi)}}{\alpha_m(\xi)}.
\]

Under the condition \( \|I - P_m(\xi)A(\xi)\|_F < 1 \), equations (3) and (4) imply that

\[
\kappa(P_m(\xi)A(\xi)) \leq \frac{\sqrt{n - (n - 1)(1 - \|I - P_m(\xi)A(\xi)\|_F^2)}}{1 - \|I - P_m(\xi)A(\xi)\|_F}.
\]
For all \( \lambda \in \mathbb{R}^m \), we have
\[
\|I - \sum_{i=1}^{m} \lambda_i P_i A(\xi)\|^2_F = n - 2\lambda^T S(\xi) + \lambda^T M(\xi) \lambda,
\]
where the matrix \( M(\xi) \in \mathbb{R}^{m \times m} \) and the vector \( S(\xi) \in \mathbb{R}^{m} \) are given by
\[
M_{i,j}(\xi) = \text{trace}(A^T(\xi) P_i^T P_j A(\xi)) \quad \text{and} \quad S_i(\xi) = \text{trace}(P_i A(\xi)).
\]

Therefore, the solution of problem (2) is \( P_m(\xi) = \sum_{i=1}^{m} \lambda_i(\xi) P_i \) with \( \lambda(\xi) \) the solution of \( M(\xi) \lambda(\xi) = S(\xi) \). When considering a small number \( m \) of interpolation points, the computation time for solving this system of equations is negligible. However, the computation of \( M(\xi) \) and \( S(\xi) \) requires the evaluation of traces of matrices \( A^T(\xi) P_i^T P_j A(\xi) \) and \( P_i A(\xi) \) for all \( 1 \leq i, j \leq m \). Since the \( P_i \) are implicit matrices, the computation of such products of matrices is not affordable for large matrices. Of course, since \( \text{trace}(B) = \sum_{i=1}^{n} e_i^T B e_i \), the trace of an implicit matrix \( B \) could be obtained by computing the product \( B \) with the canonical vectors \( e_1, \ldots, e_n \), but this approach is clearly not affordable for large \( n \).

In the next section, we propose an approximation of the above construction using an approximation of the Frobenius norm which requires less computational efforts.

### 2.2 Projection using a Frobenius semi-norm

Here, we define an approximation \( P_m(\xi) \) of \( A(\xi)^{-1} \) in \( Y_m \) by
\[
P_m(\xi) = \arg\min_{P \in Y_m} \|(I - PA(\xi))V\|_F,
\]
where \( V \in \mathbb{R}^{n \times K} \), with \( K \leq n \). \( B \mapsto \|BV\|_F \) defines a semi-norm on \( \mathbb{R}^{n \times n} \). Here, we assume that the linear map \( P \mapsto PA(\xi)V \) is injective on \( Y_m \) so that the solution of (5) is unique. This requires \( K \geq m \). Then, the solution \( P_m(\xi) = \sum_{i=1}^{m} \lambda_i(\xi) P_i \) of (5) is such that the vector \( \lambda(\xi) \in \mathbb{R}^{m} \) satisfies \( M^V(\xi) \lambda(\xi) = S^V(\xi) \), with
\[
M^V_{i,j}(\xi) = \text{trace}(V^T A^T(\xi) P_i^T P_j A(\xi)V) \quad \text{and} \quad S_i^V(\xi) = \text{trace}(V^T P_i A(\xi)V).
\]

The procedure for the computation of \( M^V(\xi) \) and \( S^V(\xi) \) is given in Algorithm 1. Note that only \( mK \) matrix-vector products involving the implicit matrices \( P_i \) are required.

Now the question is to choose a matrix \( V \) such that \( \|(I - PA(\xi))V\|_F \) provides a good approximation of \( \|I - PA(\xi)\|_F \) for any \( P \in Y_m \) and \( \xi \in \Xi \).
Algorithm 1 Computation of $M^V(\xi)$ and $S^V(\xi)$.

**Require:** $A(\xi)$, $(P_1, \ldots, P_m)$ and $V = (v_1, \ldots, v_K)$

**Ensure:** $M^V(\xi)$ and $S^V(\xi)$

1: Compute the vectors $w_{i,k} = P_i A(\xi) v_k \in \mathbb{R}^n$, for $1 \leq k \leq K$ and $1 \leq i \leq m$;
2: Set $W_i = (w_{i,1}, \ldots, w_{i,K}) \in \mathbb{R}^{n \times K}$, $1 \leq i \leq m$;
3: Compute $M_{i,j}^V(\xi) = \text{trace}(W_i^T W_j)$ for $1 \leq i, j \leq m$;
4: Compute $S_i^V(\xi) = \text{trace}(V_i^T W_i)$ for $1 \leq i \leq m$;

2.2.1 Statistical estimation of the Frobenius norm of an implicit matrix

Let $B$ an implicit $n$-by-$n$ matrix (consider $B = I - PA(\xi)$, with $P \in Y_m$ and $\xi \in \Xi$). For the computation of the Frobenius norm $\|B\|_F = \text{trace}(B^T B)^{1/2}$, we can use a statistical estimator for the trace of implicit matrices as first proposed in [22]. Let $z$ be a random vector of length $n$ whose coefficients are independent and identically distributed with zero mean and unit variance. Then $\|B\|^2_F = \mathbb{E}(\|Bz\|^2_2)$, where $\mathbb{E}(\cdot)$ denotes the mathematical expectation and $\|\cdot\|_2$ is the canonical norm of $\mathbb{R}^n$, so that $\|B\|^2_F$ can be approximated by the unbiased statistical estimator

$$T_K(B) = \frac{1}{K} \sum_{k=1}^K \|Bz_k\|_2^2,$$  \hspace{1cm} (7)

where $z_k$ are $K$ independent copies of $z$. This estimator can be written $T_K(B) = \|BV\|^2_F$, where $V$ is the random matrix whose $k$-th column is the vector $z_k/\sqrt{K}$. Different probability distributions for $z$ have been analyzed in [2]. Here, we choose the Rademacher’s distribution (entries of $z$ take values $-1$ or $1$ with probability $1/2$) which gives the lowest variance of the estimator (7). We will call $V$ the rescaled Rademacher matrix. When the matrix $B$ is diagonal, the estimator is exact for any $K \geq 1$. Theorem 13 in [2] provides a $(\varepsilon, \delta)$-confidence interval: for any $\delta \in (0, 1)$ and $\varepsilon \in (0, 1)$, if $K \geq 6 \varepsilon^{-2} \log(2n/\delta)$, then

$$(1 - \varepsilon)\|B\|^2_F \leq \|BV\|^2_F \leq (1 + \varepsilon)\|B\|^2_F$$ \hspace{1cm} (8)

holds with probability $1 - \delta$. A direct consequence of (8) is the following quasi-optimality result.

**Proposition 2.2** Let $V \in \mathbb{R}^{n \times K}$ be a random matrix such that (8) holds with a probability $1 - \delta$. Let $P_m(\xi)$ be defined by (5). Then

$$\|I - P_m(\xi)A(\xi)\|_F \leq \sqrt{\frac{1 + \varepsilon}{1 - \varepsilon}} \min_{P \in Y_m} \|I - PA(\xi)\|_F$$ \hspace{1cm} (9)

holds with probability $1 - \delta$. 

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Proof. From inequality (8) and definition (5) of $P_m(\xi)$, it holds with probability $1 - \delta$:

$$
\|I - P_m(\xi)A(\xi)\|_F \leq \frac{1}{\sqrt{1 - \varepsilon}}\|(I - P_m(\xi)A(\xi))V\|_F
$$

$$
\leq \frac{1}{\sqrt{1 - \varepsilon}} \min_{P \in Y_m} \|(I - PA(\xi))V\|_F \leq \sqrt{\frac{1 + \varepsilon}{1 - \varepsilon}} \min_{P \in Y_m} \|I - PA(\xi)\|_F \quad \blacksquare
$$

Proposition 2.2 makes the use of the statistical estimation particularly interesting for large matrices. Indeed, with $\varepsilon = 0.6$ and $\delta = 1\%$, the inequality (9) holds with $\sqrt{(1 + \varepsilon)/(1 - \varepsilon)} = 2$ and with probability 99% if $K \geq 242$ when $n = 10^4$, $K \geq 319$ when $n = 10^6$, or $K \geq 396$ when $n = 10^8$. Similarly to Proposition 2.1, we obtain the following properties for $P_m(\xi)A(\xi)$, with $P_m(\xi)$ the solution of (5).

**Proposition 2.3** Under the assumptions of Proposition 2.2, the inequalities

$$(1 - \alpha_m(\xi))^2(1 - \varepsilon)^{-1} \leq \|(I - P_m(\xi)A(\xi))V\|_F^2 \leq n (1 - (1 - \varepsilon)\alpha_m^2(\xi)) \quad (10)$$

and

$$
\kappa(P_m(\xi)A(\xi)) \leq \alpha_m(\xi)^{-1}\sqrt{n(1 - \varepsilon)^{-1} - (n - 1)\alpha_m^2(\xi)} \quad (11)
$$

hold with probability $1 - \delta$, where $\alpha_m(\xi)$ is the lowest singular value of $P_m(\xi)A(\xi)$.

Proof: The optimality condition for $P_m(\xi)$ yields $\|(I - P_m(\xi)A(\xi))V\|_F^2 = \|V\|_F^2 - \|P_m(\xi)A(\xi)V\|_F^2$. Then, using (8) and $\|V\|_F^2 = n$, we obtain that

$$
\|(I - P_m(\xi)A(\xi))V\|_F^2 \leq n - (1 - \varepsilon)\|P_m(\xi)A(\xi)\|_F^2
$$

holds with probability $1 - \delta$. Then, $\|P_m(\xi)A(\xi)\|_F^2 \geq n\alpha_m(\xi)^2$ yields the second inequality of (10). Following the proof of Lemma 2.6 in [20], we have $(1 - \alpha_m(\xi)^2) \leq \|I - P_m(\xi)A(\xi)\|_F^2$. Together with (8), it yields the first inequality of (10). Furthermore, with probability $1 - \delta$, we have $n - (1 - \varepsilon)\|P_m(\xi)A(\xi)\|_F^2 \geq 0$. Since the square of the Frobenius norm of matrix $P_m(\xi)A(\xi)$ is the sum of squares of its singular values, we deduce

$$(n - 1)\alpha_m(\xi)^2 + \beta_m(\xi)^2 \leq \|P_m(\xi)A(\xi)\|_F^2 \leq n(1 - \varepsilon)^{-1},$$

where $\beta_m(\xi)$ is the largest singular value of $P_m(\xi)A(\xi)$. Then (11) follows from the definition of $\kappa(P_m(\xi)A(\xi)) = \beta_m(\xi)/\alpha_m(\xi)$. \blacksquare
2.2.2 Use of Hadamard matrices for the estimation of the Frobenius norm of an implicit matrix

Following [3], we show how Hadamard matrices can be used for the estimation of the Frobenius norm of an implicit matrix. The goal is to find a matrix $V$ such that $\|BV\|_F$ is a good approximation of $\|B\|_F$. The relation $\|BV\|_F^2 = \text{trace}(B^TBV^T)$ suggests that $V$ should be such that $VV^T$ is as close as possible to the identity matrix. In other words, we would like $V$ to minimize

$$\text{err}(V)^2 = \frac{1}{n(n-1)} \sum_{i=1}^{n} \sum_{j \neq i} (VV^T)_{i,j}^2 = \frac{\|I - VV^T\|_F^2}{n(n-1)},$$

which is the mean square magnitude of the off-diagonal entries of $VV^T$. The bound $\text{err}(V) \geq \sqrt{(n-K)/((n-1)K)}$ holds for any $V \in \mathbb{R}^{n \times K}$ whose rows have unit norm [31]. Hadamard matrices can be used to construct matrices $V$ such that the corresponding error $\text{err}(V)$ is close to the bound.

A Hadamard matrix $H_s$ is a $s$-by-$s$ matrix whose entries are $\pm 1$, and which satisfies $H_sH_s^T = sI$ where $I$ is the identity matrix of size $s$. For example,

$$H_2 = \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

is a Hadamard matrix of size $s = 2$. The Kronecker product of two Hadamard matrices is again a Hadamard matrix. Then it is possible to build a Hadamard matrix whose size $s$ is a power of 2 using a recursive procedure: $H_{2^{k+1}} = H_{2^k} \otimes H_{2^k}$. The $(i,j)$-entry of this matrix is $(-1)^{a^Tb}$, where $a$ and $b$ are the binary vectors such that $i = \sum_{k \geq 0} 2^k a_k$ and $j = \sum_{k \geq 0} 2^k b_k$.

For a sufficiently large $s = 2^k \geq \max(n,K)$, we define the rescaled partial Hadamard matrix $V \in \mathbb{R}^{n \times K}$ as the first $n$ rows and the first $K$ columns of $H_s/\sqrt{K}$. The rows of $V$ have unit norm. We see on Figure 1(a) that the corresponding error $\text{err}(V)$ is close to the lower bound $\sqrt{(n-K)/((n-1)K)}$, and is lower than the error $\text{err}(V)$ obtained using one realization of the Rademacher random matrix introduced in Section 2.2.1. In fact, for the particular case where $n$ is a power of two, the partial Hadamard matrix reaches the bound. Furthermore, an interesting property is that $VV^T$ has a structured pattern (see Figure 1(b)). If $K = 2^q$, $VV^T$ have non-zero entries only on the $2^{qk}$-th upper and lower diagonals, with $k \geq 0$. As a consequence, the error on the estimation of $\|B\|_F$ will be induced only by the non-zero off-diagonal entries of $B$ that occupy the $2^{qk}$-th upper and lower diagonals, with $k \geq 1$. If the entries of $B$ vanish away from the diagonal, the Frobenius norm can be estimated very accurately. The rescaled Rademacher matrix can not take advantage of such particular properties of $B$ (see Figure 1(d)). However, neither the confidence interval (8) nor Proposition 2.2 hold anymore since the rescaled partial Hadamard matrix is deterministic.
We now consider the *Subsampled Randomized Hadamard Transform* (SRHT) matrix first introduced in [1]. It is defined by $K^{-1/2}RH_sD \in \mathbb{R}^{K \times s}$ where $D \in \mathbb{R}^{s \times s}$ is a diagonal random matrix where $D_{i,i}$ are independent Rademacher random variables, and $R \in \mathbb{R}^{K \times s}$ is a subset of $K$ rows from the identity matrix of size $s$ chosen uniformly at random and without replacement. In other words, we randomly select $K$ rows of $H_s$ without replacement, and we multiply the columns by $\pm K^{-1/2}$. We can find in [30, 6] an analysis of the SRHT matrix properties.

Let us now defined the partial SRHT (P-SRHT) matrix $V \in \mathbb{R}^{n \times K}$ as the first $n$ rows of $K^{-1/2}(DH_sR)^T$, with $s = 2^{\log_2(n)}$ the smallest power of 2 such that $n \leq s < 2n$. In the spirit of Lemma 4.10 in [6], we give a confidence interval of the estimator $\|BV\|_F$.

**Proposition 2.4** Let $B \in \mathbb{R}^{n \times n}$ and $V \in \mathbb{R}^{n \times K}$ a P-SRHT matrix corresponding to the first $n$ rows of a SRHT matrix $K^{-1/2}(DH_sR)^T \in \mathbb{R}^{s \times K}$ with $s = 2^{\log_2(n)}$. For any $\delta \in (0, 1)$ and $\varepsilon \in (0, 1)$, if

$$K \geq 2(\varepsilon^2 - \varepsilon^3/3)^{-1} \ln(2/\delta)(1 + \sqrt{8\ln(2n/\delta)})^2,$$

then (8) holds with a probability of $1 - 2\delta$.

**Proof:** We define the square matrix $\tilde{B}$ of size $s$, whose first $n \times n$ diagonal block is $B$, and 0 elsewhere. Then we have $\|\tilde{B}\|_F = \|B\|_F$. The rest of the proof is similar to the one of Lemma 4.10 in [6]. Let $A = \{(1 - \varepsilon)\|\tilde{B}\|_F^2 \leq \|\tilde{B}V\|_F^2 \leq (1 + \varepsilon)\|\tilde{B}\|_F^2\}$ and $E = \{\sup_i \|\tilde{B}DH_se_i\|_F^2 \leq (1 + \sqrt{8\ln(s/\delta)})^2\|\tilde{B}\|_F^2\}$. The relation $\mathbb{P}(A^c) \leq \mathbb{P}(A^c|E) + \mathbb{P}(E^c)$ holds. Thanks to Lemma 4.6 in [6] we have $\mathbb{P}(E^c) \leq \delta$. Using the scalar Chernoff bound (Theorem 2.2 in [30] with $k = 1$) we have

$$\mathbb{P}(A^c|E) = \mathbb{P}(\|\tilde{B}V\|_F^2 \leq (1 - \varepsilon)\|\tilde{B}\|_F^2|E) + \mathbb{P}(\|\tilde{B}V\|_F^2 \geq (1 + \varepsilon)\|\tilde{B}\|_F^2|E) \leq (e^{\varepsilon}(1 + \varepsilon)^{-1} - \varepsilon)^K(1 + \sqrt{8\ln(s/\delta)})^{-2} + (e^{-\varepsilon}(1 - \varepsilon)^{-1} + \varepsilon)^K(1 + \sqrt{8\ln(s/\delta)})^{-2} \leq 2(e^{\varepsilon}(1 + \varepsilon)^{-1} - \varepsilon)^K(1 + \sqrt{8\ln(s/\delta)})^{-2} \leq 2e^{K(-\varepsilon^2/2 + 3\varepsilon^3/6)(1 + \sqrt{8\ln(s/\delta)})^{-2}}.$$

The condition (12) implies $\mathbb{P}(A^c) \leq 2\delta$, which ends the proof.

Figure 1(a) shows that the error $err(V)$ associated to the P-SRHT matrix behaves almost like the rescaled partial Hadamard matrix. Furthermore the confidence interval given in proposition 2.4 is similar to the one associated to the Rademacher matrix: according to (12) $K$ scales as $O(\varepsilon^{-2} \log(\delta^{-1})^2 \log(n))$ when $n \to \infty$ and $(\delta, \varepsilon) \to 0$. Up to a factor $\log(\delta^{-1})$, this is the same behavior as for the Rademacher matrix.
Figure 1: Comparison between the rescaled Rademacher, the rescaled partial Hadamard and the P-SRHT matrix for the definition of matrix $V$, with $n = 600$.

### 2.3 Ensuring the invertibility of the preconditioner for positive definite matrix

Here, we propose a modification of the interpolation which ensures that $P_m(\xi)$ is invertible when $A(\xi)$ is positive definite.

Since $A(\xi_i)$ is positive definite, $P_i = A(\xi_i)^{-1}$ is positive definite. We introduce the vectors
\( \gamma^- \in \mathbb{R}^m \) and \( \gamma^+ \in \mathbb{R}^m \) whose components
\[
\gamma^-_i = \inf_{w \in \mathbb{R}^n} \frac{\langle P_i w, w \rangle}{\|w\|^2} > 0 \quad \text{and} \quad \gamma^+_i = \sup_{w \in \mathbb{R}^n} \frac{\langle P_i w, w \rangle}{\|w\|^2} < \infty
\]
correspond respectively to the lowest and highest eigenvalues of the symmetric part of \( P_i \).

Then, for any \( P = \sum_{i=1}^m \lambda_i P_i \in Y_m \),
\[
\inf_{w \in \mathbb{R}^n} \frac{\langle P w, w \rangle}{\|w\|^2} \geq \langle \lambda^+, \gamma^- \rangle - \langle \lambda^-, \gamma^+ \rangle,
\]
where \( \lambda^+ \geq 0 \) and \( \lambda^- \geq 0 \) are respectively the positive and negative parts of \( \lambda = \lambda^+ - \lambda^- \in \mathbb{R}^m \). As a consequence, if the right hand side of (13) is strictly positive, then \( P \) is invertible. Furthermore, we have \( \|P\| \leq \frac{\langle \lambda^+ + \lambda^-, C \rangle}{\langle \lambda^+, \gamma^- \rangle - \langle \lambda^-, \gamma^+ \rangle} \), where \( C \in \mathbb{R}^m \) is the vector of component \( C_i = \|P_i\| \), where \( \|P_i\| \) denotes the operator norm of \( P_i \). If we assume that \( \langle \lambda^+, \gamma^- \rangle - \langle \lambda^-, \gamma^+ \rangle > 0 \), the condition number of \( P \) satisfies
\[
\kappa(P) = \frac{\|P\| \|P^{-1}\|}{\inf_{w \in \mathbb{R}^n} \frac{\langle P w, w \rangle}{\|w\|^2}} \leq \frac{\langle \lambda^+ + \lambda^-, C \rangle}{\langle \lambda^+, \gamma^- \rangle - \langle \lambda^-, \gamma^+ \rangle}.
\]

It is then possible to bound \( \kappa(P) \) by \( \bar{\kappa} \) by imposing
\[
\langle \lambda^+ + \lambda^-, C \rangle \leq \bar{\kappa}(\langle \lambda^+, \gamma^- \rangle - \langle \lambda^-, \gamma^+ \rangle),
\]
which is a linear inequality constraint on \( \lambda^+ \) and \( \lambda^- \). We introduce two convex subsets of \( Y_m \) defined by
\[
Y^\bar{\kappa}_m = \left\{ \sum_{i=1}^m \lambda_i^+ P_i - \sum_{i=1}^m \lambda_i^- P_i : \lambda_i^+ \geq 0, \lambda_i^- \geq 0, \langle \lambda^+, \gamma^- \rangle - \langle \lambda^-, \gamma^+ \rangle \geq 0 \right\},
\]
\[
Y^+_m = \left\{ \sum_{i=1}^m \lambda_i P_i : \lambda_i \geq 0 \right\}.
\]

From (13), we have that any nonzero element of \( Y^+_m \) is invertible, while any nonzero element of \( Y^\bar{\kappa}_m \) is invertible and has a condition number lower than \( \bar{\kappa} \). Under the condition \( \bar{\kappa} \geq \max_i C_i/\gamma^-_i \), we have
\[
Y^+_m \subset Y^\bar{\kappa}_m \subset Y_m.
\]

Then definitions (2) and (5) for the approximation \( P_m(\xi) \) can be replaced respectively by
\[
P_m(\xi) = \arg\min_{P \in Y^+_m \text{ or } Y^\bar{\kappa}_m} \|I - PA(\xi)\|_F, \quad (16a)
\]
\[
P_m(\xi) = \arg\min_{P \in Y^+_m \text{ or } Y^\bar{\kappa}_m} \|(I - PA(\xi))V\|_F, \quad (16b)
\]
which are quadratic optimization problems with linear inequality constraints. Furthermore, since $P_i \in Y_+^m$ for all $i$, all the resulting projections $P_m(\xi)$ interpolate $A(\xi)^{-1}$ at the points $\xi_1, \ldots, \xi_m$.

The following proposition shows that properties (3) and (4) still hold for the preconditioned operator.

**Proposition 2.5** The solution $P_m(\xi)$ of (16a) is such that $P_m(\xi)A(\xi)$ satisfies (3) and (4). Also, the solution $P_m(\xi)$ of (16b) is such that $P_m(\xi)A(\xi)$ satisfies (10) and (11).

**Proof:** Since $Y_+^m$ (or $Y^\kappa_m$) is a closed and convex positive cone, the solution $P_m(\xi)$ of (16a) is such that $\text{trace}((I - P_m(\xi)A(\xi))^T(P_m(\xi) - P)A(\xi)) \geq 0$ for all $P \in Y_+^m$ (or $Y^\kappa_m$). Taking $P = 2P_m(\xi)$ and $P = 0$, we obtain that $\text{trace}((I - P_m(\xi)A(\xi))^TP_m(\xi)A(\xi)) = 0$, which implies $\|P_m(\xi)A(\xi)\|_F^2 = \text{trace}(P_m(\xi)A(\xi))$. We refer to the proof of Lemma 2.6 and Theorem 3.2 in [20] to deduce (3) and (4). Using the same arguments, we prove that the solution $P_m(\xi)$ of (16b) satisfies $\|P_m(\xi)A(\xi)V\|_F^2 = \text{trace}(V^TP_m(\xi)A(\xi)V)$, and then that (10) and (11) hold. ■

### 2.4 Practical computation of the projection

Here, we detail how to efficiently compute $M^V(\xi)$ and $S^V(\xi)$ given in equation (6) in a multi-query context, i.e. for several different values of $\xi$. The same methodology can be applied for computing $M(\xi)$ and $S(\xi)$. We assume that the operator $A(\xi)$ has an affine expansion of the form

$$A(\xi) = \sum_{k=1}^{m_A} \Phi_k(\xi) A_k,$$

where the $A_k$ are matrices in $\mathbb{R}^{n \times n}$ and the $\Phi_k : \Xi \rightarrow \mathbb{R}$ are real valued functions. Then $M^V(\xi)$ and $S^V(\xi)$ also have the affine expansions

$$M^V_{i,j}(\xi) = \sum_{k=1}^{m_A} \sum_{l=1}^{m_A} \Phi_k(\xi) \Phi_l(\xi) \text{trace}(V^T A^T_k P^T_i P^T_j A_l V),$$

$$S^V_i(\xi) = \sum_{k=1}^{m_A} \Phi_k(\xi) \text{trace}(V^T P_i A_k V),$$

respectively. Computing the multiple terms of these expansions would require many computations of traces of implicit matrices and also, it would require the computation of the affine expansion of $A(\xi)$. Here, we use the methodology introduced in [8] for obtaining affine
decompositions with a lower number of terms. These decompositions only require the knowledge
of functions $\Phi_k$ in the affine decomposition \cite{17}, and evaluations of $M_{ij}(\xi)$ and $S_{ij}(\xi)$
(that means evaluations of $A(\xi)$) at some selected points. We briefly recall this methodology.

Suppose that $g : \Xi \to X$, with $X$ a vector space, has an affine decomposition $g(\xi) = \sum_{k=1}^{m} \zeta_k(\xi) g_k$, with $\zeta_k : \Xi \to \mathbb{R}$ and $g_k \in X$. We first compute an interpolation of $\zeta(\xi) = (\xi_1(\xi), \ldots, \xi_m(\xi))$ under the form $\zeta(\xi) = \sum_{k=1}^{m} \Psi_k(\xi) \zeta_k(\xi)$, with $m_{g_1} \leq m$, where $\xi_1^*, \ldots, \xi_m^*$ are interpolation points and $\Psi_1(\xi), \ldots, \Psi_m(\xi)$ the associated interpolation functions. Such an interpolation can be computed with the Empirical Interpolation Method \cite{25} described in Algorithm 2. Then, we obtain an affine decomposition $g(\xi) = \sum_{k=1}^{m_{\Psi}} \Psi_k(\xi) g(\xi_k^*)$ which can be computed from evaluations of $g$ at interpolation points $\xi_k^*$.

**Algorithm 2:** Empirical Interpolation Method (EIM).

---

**Require:** $(\xi_1(\cdot), \ldots, \xi_m(\cdot))$

**Ensure:** $\Psi_1(\cdot), \ldots, \Psi_m(\cdot)$ and $\xi_1^*, \ldots, \xi_m^*$

1: Define $R_1(i, \xi) = \zeta_i(\xi)$ for all $i, \xi$
2: Initialize $e = 1$, $k = 0$
3: while $e \geq tolerance$ (in practice the machine precision) do
4:   $k = k + 1$
5:   Find the pivot $(i_k^*, \xi_k^*) \in \text{argmax}_{i, \xi} |R_k(i, \xi)|$
6:   Set the error to $e = |R_k(i_k^*, \xi_k^*)|$
7:   Actualize $R_{k+1}(i, \xi) = R_k(i, \xi) - R_k(i, \xi_k^*) R_k(i_k^*, \xi) / R_k(i_k^*, \xi_k^*)$ for all $i, \xi$
8: end while
9: Fill in the $k$-by-$k$ matrix $Q : Q_{ij} = \zeta_i^*(\xi_j^*)$ for all $1 \leq i, j \leq k$
10: Compute $\Psi_i(\xi) = \sum_{j=1}^{k} (Q^{-1})_{ij} \zeta_j(\xi)$ for all $i, \xi$

---

Applying the above procedure to both $M(\xi)$ and $S(\xi)$, we obtain

$$M(\xi) \approx \sum_{k=1}^{m_{M}} \Psi_k(\xi) M(\xi_k^*), \quad S(\xi) \approx \sum_{k=1}^{m_{S}} \Psi_k(\xi) S(\xi_k^*).$$

(19)

The first (so called offline) step consists in computing the interpolation functions $\Psi_k(\xi)$ and $\Psi_k(\xi)$ and associated interpolation points $\xi_k^*$ and $\xi_k^*$ using Algorithm 2 with input $\{\Phi \Phi\}_{1 \leq i, j \leq m_A}$ and $\{\Phi_1\}_{1 \leq i \leq m_A}$ respectively, and then in computing matrices $M(\xi_k^*)$ and vectors $S(\xi_k^*)$ using Algorithm 1. The second (so called online) step simply consists in computing the matrix $M(\xi)$ and the vector $S(\xi)$ for a given value of $\xi$ using (19).
3 Preconditioners for projection-based model reduction

We consider a parameter-dependent linear equation

$$A(\xi)u(\xi) = b(\xi),$$

with $A(\xi) \in \mathbb{R}^{n \times n}$ and $b(\xi) = \mathbb{R}^n$. Projection-based model reduction consists in projecting the solution $u(\xi)$ onto an approximation space $X_r \subset X := \mathbb{R}^n$ of low dimension $r \ll n$. In this section, we show how the preconditioner $P_m(\xi)$ can be used for the definition of the projection and for the construction of the approximation space $X_r$.

$X$ is endowed with the norm $\| \cdot \|_X$ defined by $\| \cdot \|_X^2 = \langle R_X \cdot, \cdot \rangle$, where $R_X$ is a symmetric positive definite matrix and $\langle \cdot, \cdot \rangle$ is the canonical inner product of $\mathbb{R}^n$. We also introduce the dual norm $\| \cdot \|_{X^\prime} = \| R_X^{-1} \cdot \|_X$ such that for any $v, w \in X$ we have $|\langle v, w \rangle| \leq \|v\|_X \|w\|_{X^\prime}$.

3.1 Projection of the solution on a given reduced subspace

Here, we suppose that the approximation space $X_r$ has been computed by some model order reduction method. The best approximation of $u(\xi)$ on $X_r$ is $u^*_r(\xi) = \arg \min_{v \in X_r} \|u(\xi) - v\|_X$ and is characterized by the orthogonality condition

$$\langle u^*_r(\xi) - u(\xi), R_X v_r \rangle = 0, \quad \forall v_r \in X_r,$$

or equivalently by the Petrov-Galerkin orthogonality condition

$$\langle A(\xi) u^*_r(\xi) - b(\xi), A^{-T}(\xi) R_X v_r \rangle = 0, \quad \forall v_r \in X_r. \quad (21)$$

Obviously the computation of test functions $A^{-T}(\xi) R_X v_r$ for basis functions $v_r$ of $X_r$ is prohibitive. By replacing $A(\xi)^{-1}$ by $P_m(\xi)$, we obtain the feasible Petrov-Galerkin formulation

$$\langle A(\xi) u_r(\xi) - b(\xi), P_m^T(\xi) R_X v_r \rangle = 0, \quad \forall v_r \in X_r. \quad (22)$$

Denoting by $U \in \mathbb{R}^{n \times r}$ a matrix whose range is $X_r$, the solution of (23) is $u_r(\xi) = U a(\xi)$ where the vector $a(\xi) \in \mathbb{R}^r$ is the solution of

$$(U^T R_X P_m(\xi) A(\xi) U) a(\xi) = U^T R_X P_m(\xi) b(\xi).$$

Note that (23) corresponds to the standard Galerkin projection when replacing $P_m(\xi)$ by $R_X^{-1}$.

We give now a quasi-optimality result for the approximation $u_r(\xi)$. This analysis relies on the notion of $\delta$-proximality introduced in [12].
Proposition 3.1 Let $\delta_{r,m}(\xi) \in [0,1]$ be defined by

$$\delta_{r,m}(\xi) = \max_{v_r \in X_r} \min_{w_r \in X_r} \frac{\|v_r - R^{-1}_X(P_m(\xi)A(\xi))^T R_X w_r\|_X}{\|v_r\|_X}. \quad (24)$$

The solutions $u^*_r(\xi) \in X_r$ and $u_r(\xi) \in X_r$ of (21) and (23) satisfy

$$\|u^*_r(\xi) - u_r(\xi)\|_X \leq \delta_{r,m}(\xi)\|u(\xi) - u_r(\xi)\|_X. \quad (25)$$

Moreover, if $\delta_{r,m}(\xi) < 1$ holds, then

$$\|u(\xi) - u_r(\xi)\|_X \leq (1 - \delta_{r,m}(\xi)^2)^{-1/2}\|u(\xi) - u^*_r(\xi)\|_X. \quad (26)$$

Proof: The orthogonality condition (21) yields

$$\langle u^*_r(\xi) - u_r(\xi), R_X v_r \rangle = \langle u(\xi) - u_r(\xi), R_X v_r \rangle = \langle b(\xi) - A(\xi)u_r(\xi), A^{-T}(\xi) R_X v_r \rangle$$

for all $v_r \in X_r$. Using (23), we have that for any $w_r \in X_r$,

$$\langle u^*_r(\xi) - u_r(\xi), R_X v_r \rangle = \langle b(\xi) - A u_r(\xi), A^{-T}(\xi) R_X v_r - P_m(\xi)^T R_X w_r \rangle,
= \langle u(\xi) - u_r(\xi), R_X v_r - (P_m(\xi)A(\xi))^T R_X w_r \rangle,
\leq \|u(\xi) - u_r(\xi)\|_X \|R_X v_r - (P_m(\xi)A(\xi))^T R_X w_r\|_X,
= \|u(\xi) - u_r(\xi)\|_X \|v_r - R^{-1}_X(P_m(\xi)A(\xi))^T R_X w_r\|_X.$$

Taking the infimum over $w_r \in X_r$ and by the definition of $\delta_{r,m}(\xi)$, we obtain

$$\langle u^*_r(\xi) - u_r(\xi), R_X v_r \rangle \leq \delta_{r,m}(\xi)\|u(\xi) - u_r(\xi)\|_X \|v_r\|_X.$$

Then, noting that $u^*_r(\xi) - u_r(\xi) \in X_r$, we obtain

$$\|u^*_r(\xi) - u_r(\xi)\|_X = \sup_{v_r \in X_r} \frac{\langle u^*_r(\xi) - u_r(\xi), R_X v_r \rangle}{\|v_r\|_X} \leq \delta_{r,m}(\xi)\|u(\xi) - u_r(\xi)\|_X,$$

that is (25). Finally, using orthogonality condition (21), we have that

$$\|u(\xi) - u_r(\xi)\|^2_X = \|u(\xi) - u^*_r(\xi)\|^2_X + \|u^*_r(\xi) - u_r(\xi)\|^2_X,
\leq \|u(\xi) - u^*_r(\xi)\|^2_X + \delta_{r,m}(\xi)^2\|u(\xi) - u_r(\xi)\|^2_X,$$

from which we deduce (26) when $\delta_{r,m}(\xi) < 1$. \hfill \blackslug

An immediate consequence of Proposition 3.1 is that when $\delta_{r,m}(\xi) = 0$, the Petrov-Galerkin projection $u_r(\xi)$ coincides with the orthogonal projection $u^*_r(\xi)$. Following [11], we show in the following proposition that $\delta_{r,m}(\xi)$ can be computed by solving an eigenvalue problem of size $r$. 

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Proposition 3.2 We have \( \delta_{r,m}(\xi) = \sqrt{1 - \gamma} \), where \( \gamma \) is the lowest eigenvalue of the generalized eigenvalue problem \( Cx = \gamma Dx \), with
\[
C = U^T B (B^T R_X^{-1} B)^{-1} B^T U \in \mathbb{R}^{r \times r},
\]
\[
D = U^T R_X U \in \mathbb{R}^{r \times r},
\]
where \( B = (P_m(\xi) A(\xi))^T R_X U \in \mathbb{R}^{n \times r} \) and \( U \in \mathbb{R}^{n \times r} \) is a matrix whose range is \( X_r \).

Proof: Since the range of \( U \) is \( X_r \), we have
\[
\delta_{r,m}(\xi)^2 = \max_{a \in \mathbb{R}^r \setminus \{0\}} \min_{b \in \mathbb{R}^r} \frac{\|Ua - R_X^{-1} B b\|_X^2}{\|Ua\|_X^2}. 
\]

For any \( a \in \mathbb{R}^r \), the minimizer \( b^* \) of \( \|Ua - R_X^{-1} B b\|_X^2 \) over \( b \in \mathbb{R}^r \) is given by \( b^* = (B^T R_X^{-1} B)^{-1} B^T U a \). Therefore, we have
\[
\delta_{r,m}(\xi)^2 = 1 - \inf_{a \in \mathbb{R}^r} \frac{\langle U^T B (B^T R_X^{-1} B)^{-1} B^T U a, a \rangle}{\langle U^T R_X U a, a \rangle}, 
\]
which concludes the proof.

3.2 Greedy construction of the solution reduced subspace

Following the idea of the Reduced Basis method \cite{29}, a sequence of nested approximation spaces \( \{X_r\}_{r \geq 1} \) in \( X \) can be constructed by a greedy algorithm such that \( X_{r+1} = X_r + \text{span}(u(\xi^{RB}_{r+1})) \), where \( \xi^{RB}_{r+1} \) is a point where the error of approximation of \( u(\xi) \) in \( X_r \) is maximal. An ideal greedy algorithm using the best approximation in \( X_r \) and an exact evaluation of the projection error is such that
\[
u_r^*(\xi) \text{ is the orthogonal projection of } u(\xi) \text{ on } X_r \text{ defined by } (21),
\]
\[
\xi^{RB}_{r+1} \in \arg\max_{\xi \in \Xi} \|u(\xi) - u_r^*(\xi)\|_X. \tag{27a}
\]

This ideal greedy algorithm is not feasible in practice since \( u(\xi) \) is not known. Therefore, we rather rely on a feasible weak greedy algorithm such that
\[
u_r(\xi) \text{ is the Petrov-Galerkin projection of } u(\xi) \text{ on } X_r \text{ defined by } (23),
\]
\[
\xi^{RB}_{r+1} \in \arg\max_{\xi \in \Xi} \|P_m(\xi) (A(\xi) u_r(\xi) - b(\xi))\|_X. \tag{28a}
\]

Assume that
\[
\alpha_m \|u(\xi) - u_r(\xi)\|_X \leq \|P_m(\xi) (A(\xi) u_r(\xi) - b(\xi))\|_X \leq \beta_m \|u(\xi) - u_r(\xi)\|_X
\]

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holds with \( \alpha_m = \inf_{\xi \in \Xi} \alpha_m(\xi) \geq 0 \) and \( \beta_m = \sup_{\xi \in \Xi} \beta_m(\xi) < \infty \), where \( \alpha_m(\xi) \) and \( \beta_m(\xi) \) are respectively the lowest and largest singular values of \( P_m(\xi)A(\xi) \) with respect to norm \( \| \cdot \|_X \), respectively defined by the infimum and supremum of \( \| P_m(\xi)A(\xi)v \|_X \) over \( v \in X \) such that \( \| v \|_X = 1 \). Then, we easily prove that algorithm (28) is such that

\[
\| u(\xi^{RB}_{r+1}) - u_r(\xi^{RB}_{r+1}) \|_X \geq \gamma_m \max_{\xi \in \Xi} \| u(\xi) - u_r(\xi) \|_X,
\]

where \( \gamma_m = \frac{\alpha_m}{\beta_m} \leq 1 \) measures how far the selection of the new point is from the ideal greedy selection. Under condition (29), convergence results for this weak greedy algorithm can be found in [5, 15].

We give now sharper bounds for the preconditioned residual norm that exploits the fact that the approximation \( u_r(\xi) \) is the Petrov-Galerkin projection.

**Proposition 3.3** Let \( u_r(\xi) \) be the Petrov-Galerkin projection of \( u(\xi) \) on \( X_r \) defined by (22). Then we have

\[
\alpha_{r,m}(\xi) \| u(\xi) - u_r(\xi) \|_X \leq \| P_m(\xi)(A(\xi)u_r(\xi) - b(\xi)) \|_X \leq \beta_{r,m}(\xi) \| u(\xi) - u_r(\xi) \|_X,
\]

with

\[
\alpha_{r,m}(\xi) = \inf_{v \in X} \sup_{w_r \in X_r} \frac{\| (P_m(\xi)A(\xi))^T R_X v \|_{X'}}{\| v - w_r \|_X},
\]

\[
\beta_{r,m}(\xi) = \sup_{v \in X} \inf_{w_r \in X_r} \frac{\| (P_m(\xi)A(\xi))^T R_X (v - w_r) \|_{X'}}{\| v \|_X}.
\]

**Proof:** For any \( v \in X \) and \( w_r \in X_r \) and according to (23), we have

\[
\langle u(\xi) - u_r(\xi), R_X v \rangle = \langle b(\xi) - A(\xi)u_r(\xi), A^{-T}(\xi)R_X v - P_m^T(\xi)R_X w_r \rangle
\]

\[
= \langle P_m(\xi)(b(\xi) - A(\xi)u_r(\xi)), (P_m(\xi)A(\xi))^{-T}R_X v - R_X w_r \rangle
\]

\[
\leq \| R \|_X \| (P_m(\xi)A(\xi))^{-T}R_X v - R_X w_r \|_{X'},
\]

where \( R(\xi) := P_m(\xi)(b(\xi) - A(\xi)u_r(\xi)) \). Taking the infimum over \( w_r \in X_r \), dividing by \( \| v \|_X \) and taking the supremum over \( v \in X \), we obtain

\[
\| u(\xi) - u_r(\xi) \|_X \leq \| R(\xi) \|_X \sup_{w_r \in X_r} \inf_{v \in X} \frac{\| (P_m(\xi)A(\xi))^{-T}R_X v - R_X w_r \|_{X'}}{\| v \|_X},
\]

\[
= \| R(\xi) \|_X \sup_{w_r \in X_r} \inf_{v \in X} \frac{\| v - w_r \|_X}{\| (P_m(\xi)A(\xi))^T R_X v \|_{X'}}.
\]

\[
= \| R(\xi) \|_X \left( \inf_{w_r \in X_r} \sup_{v \in X} \frac{\| (P_m(\xi)A(\xi))^T R_X v \|_{X'}}{\| v - w_r \|_X} \right)^{-1},
\]

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which proves the first inequality. Furthermore, for any \( v \in X \) and \( w_r \in X_r \), we have
\[
\langle P_m(\xi)(b(\xi) - A(\xi)u_r(\xi)), R_X v \rangle = \langle b(\xi) - A(\xi)u_r(\xi), P_m^T(\xi)R_X(v - w_r) \rangle 
\leq \|u(\xi) - u_r(\xi)\|_X \|(P_m(\xi)A(\xi))^T R_X(v - w_r)\|_{X'}.
\]
Taking the infimum over \( w_r \in X_r \), dividing by \( \|v\|_X \) and taking the supremum over \( v \in X \), we obtain the second inequality.

Since \( X_r \subset X_{r+1} \), we have \( \alpha_{r+1,m}(\xi) \geq \alpha_{r,m}(\xi) \geq \alpha_m(\xi) \) and \( \beta_{r+1,m}(\xi) \leq \beta_{r,m}(\xi) \leq \beta_m(\xi) \), meaning that \( \|P_m(\xi)(A(\xi)u_r(\xi) - b(\xi))\|_X \) is getting closer and closer from the true error \( \|u(\xi) - u_r(\xi)\|_X \). Equation (29) holds with \( \gamma_m \) replaced by the parameter \( \gamma_{r,m} = \alpha_{r,m}/\beta_{r,m} \). Since \( \gamma_{r,m} \) increases with \( r \), a reasonable expectation is that the convergence properties of the weak greedy algorithm will improve when \( r \) increases.

**Remark 3.4** When replacing \( P_m(\xi) \) by \( R_X^{-1} \), the preconditioned residual norm \( \|P_m(\xi)(A(\xi)u_r(\xi) - b(\xi))\|_X \) turns out to be the residual norm \( \|A(\xi)u_r(\xi) - b(\xi)\|_{X'} \), which is a standard choice in the Reduced Basis method for the greedy selection of points (with \( R_X \) being associated with the natural norm on \( X \) or with a norm associated with the operator at some nominal parameter value). This can be interpreted as a basic preconditioning method with a parameter-independent preconditioner.

## 4 Selection of the interpolation points

In this section, we propose strategies for the adaptive selection of the interpolation points. For a given set of interpolation points \( \xi_1, \ldots, \xi_m \), three different methods are proposed for the selection of a new interpolation point \( \xi_{m+1} \). The first method aims at reducing uniformly the error between the inverse operator and its interpolation. The resulting interpolation of the inverse is pertinent for preconditioning iterative solvers or estimating errors based on preconditioned residuals. The second method aims at improving Petrov-Galerkin projections of the solution of a parameter-dependent equation on a given approximation space. The third method aims at reducing the cost for the computation of the preconditioner by reusing operators computed when solving samples of a parameter-dependent equation.

### 4.1 Greedy approximation of the inverse of a parameter-dependent matrix

A natural idea is to select a new interpolation point where the preconditioner \( P_m(\xi) \) is not a good approximation of \( A(\xi)^{-1} \). Obviously, an ideal strategy for preconditioning would be
to choose $\xi_{m+1}$ where the condition number of $P_m(\xi)A(\xi)$ is maximal. The computation of
the condition number for many values of $\xi$ being computationally expensive, one could use
upper bounds of this condition number, e.g. computed using SCM [23].

Here, we propose the following selection method: given an approximation $P_m(\xi)$ associated with interpolation points $\xi_1, \ldots, \xi_m$, a new point $\xi_{m+1}$ is selected such that

$$\xi_{m+1} \in \arg\max_{\xi \in \Xi} \| (I - P_m(\xi)A(\xi)) V \|_F,$$

where the matrix $V$ is either the random Rademacher matrix, or the partial (randomized) Hadamard matrix (see Section 2.2). This adaptive selection of the interpolation points yields the construction of a growing sequence of subspaces $Y_{m+1} = Y_m + \text{span}(A(\xi_{m+1})^{-1})$ in $Y = \mathbb{R}^{n \times n}$. This algorithm is detailed below.

**Algorithm 3** Greedy selection of interpolation points.

**Require:** $A(\xi), V, M$.

**Ensure:** Interpolation points $\xi_1, \ldots, \xi_M$ and interpolation $P_M(\xi)$.

1. Initialize $P_0(\xi) = I$
2. for $m = 0$ to $M - 1$ do
3. Compute the new point $\xi_{m+1}$
4. Compute a factorization of $A(\xi_{m+1})$
5. Define $A(\xi_{m+1})^{-1}$ as an implicit operator
6. Update the space $Y_{m+1} = Y_m + \text{span}(A(\xi_{m+1})^{-1})$
7. Compute $P_{m+1}(\xi) = \arg\min_{P \in Y_{m+1}} \| (I - P A(\xi)) V \|_F$
8. end for

The following lemma interprets the above construction as a weak greedy algorithm.

**Lemma 4.1** Assume that $A(\xi)$ satisfies $\alpha_0 \| \cdot \| \leq \| A(\xi) \cdot \| \leq \bar{\beta}_0 \| \cdot \|$ for all $\xi \in \Xi$. Let $P_m(\xi)$ be defined by (5). Under the assumptions of Proposition 2.2

$$\| P_m(\xi_{m+1}) - A(\xi_{m+1})^{-1} \|_F \geq \gamma_e \max_{\xi \Xi} \min_{P \in Y_m} \| P - A(\xi)^{-1} \|_F$$

holds with probability $1 - \delta$, with $\gamma_e = \frac{\alpha_0 \gamma}{\beta_0 \sqrt{1 + \varepsilon}}$.

**Proof:** Since $\| B C \|_F \leq \| B \|_F \| C \|$ holds for any matrices $B$ and $C$, with $\| C \|$ the operator norm of $C$, we have for all $P \in Y$,

$$\| A(\xi)^{-1} - P \|_F \leq \| I - P A(\xi) \|_F \| A(\xi)^{-1} \| \leq \alpha_0^{-1} \| I - P A(\xi) \|_F,$$

$$\| I - P A(\xi) \|_F \leq \| A(\xi)^{-1} - P \|_F \| A(\xi) \| \leq \bar{\beta}_0 \| A(\xi)^{-1} - P \|_F.$$
Then, under the assumptions of Proposition 2.2, we have with probability $1 - \delta$

\[ \|A(\xi)^{-1} - P\|_F \leq (\alpha_0 \sqrt{1 - \varepsilon})^{-1} \| (I - PA(\xi)) V \|_F \]

and \[ \| (I - PA(\xi)) V \|_F \leq \beta_0 \sqrt{1 + \varepsilon} \| A(\xi)^{-1} - P \|_F, \]

which implies

\[ \frac{1}{\beta_0 \sqrt{1 + \varepsilon}} \| (I - PA(\xi)) V \|_F \leq \| A(\xi)^{-1} - P \|_F \leq \frac{1}{\alpha_0 \sqrt{1 - \varepsilon}} \| (I - PA(\xi)) V \|_F. \]

We easily deduce that $\xi_{m+1}$ is such that (32) holds with probability $1 - \delta$.

The quality of the resulting spaces $Y_m$ have to be compared with the Kolmogorov $m$-width of the set $A^{-1}(\Xi) := \{ A(\xi)^{-1} : \xi \in \Xi \} \subset Y$, defined by

\[ d_m(A^{-1}(\Xi))_Y = \min_{Y_m \subset Y} \sup_{\xi \in \Xi} \min_{P \in Y_m} \| A(\xi)^{-1} - P \|_F, \]

(33)

which evaluates how well the elements of $A^{-1}(\Xi)$ can be approximated on a $m$-dimensional subspace of matrices. Lemma 4.1 implies that the following results holds with probability $1 - \delta$ (see Corollary 3.3 in [15]):

\[ \| A(\xi)^{-1} - P_m(\xi) \|_F = \begin{cases} \mathcal{O}(m^{-a}) & \text{if } d_m(A^{-1}(\Xi))_Y = \mathcal{O}(m^{-a}) \\ \mathcal{O}(e^{-\tilde{c}m^b}) & \text{if } d_m(A^{-1}(\Xi))_Y = \mathcal{O}(e^{-\tilde{c}m^b}) \end{cases}, \]

where $\tilde{c} > 0$ is a constant which depends on $c$ and $b$. That means that if the Kolmogorov $m$-width has an algebraic or exponential convergence rate, then the weak greedy algorithm yields an error $\|P_m(\xi) - A(\xi)^{-1}\|_F$ which has the same type of convergence, with probability $1 - \delta$. Therefore, the proposed interpolation method will present good convergence properties when $d_m(A^{-1}(\Xi))_Y$ rapidly decreases with $m$.

**Remark 4.2** When the parameter set $\Xi$ is $[-1, 1]^d$ (or a product of compact intervals), an exponential decay can be obtained when $A(\xi)^{-1}$ admits an holomorphic extension to a domain in $\mathbb{C}^d$ containing $\Xi$ (see [9]). This is the case when $A(\xi)$ admits an invertible holomorphic extension to a domain in $\mathbb{C}^d$ containing $\Xi$.

**Remark 4.3** Note that here, there is no constraint on the minimization problem over $Y_m$ (either optimal subspaces or subspaces constructed by the greedy procedure), so that we have no guaranty that the resulting approximations $Y_m$ are invertible (see Section 2.3).
4.2 Selection of points for improving the projection on a reduced space

We here suppose that we want to find an approximation of the solution $u(\xi)$ of a parameter-dependent equation (20) onto a low-dimensional approximation space $X_r$, using a Petrov-Galerkin orthogonality condition given by (23). The best approximation is considered as the orthogonal projection defined by (21). The quantity $\delta_{r,m}(\xi)$ defined by (24) controls the quality of the Petrov-Galerkin projection on $X_r$ (see Proposition 3.1). As indicated in Proposition 3.2, $\delta_{r,m}(\xi)$ can be efficiently computed. Thus, we propose the following selection strategy which aims at improving the quality of the Petrov-Galerkin projection: given a preconditioner $P_m(\xi)$ associated with interpolation points $\xi_1, \ldots, \xi_m$, the next point $\xi_{m+1}$ is selected such that

$$
\xi_{m+1} \in \arg\max_{\xi \in \Xi} \delta_{r,m}(\xi).
$$

The resulting construction is described by Algorithm 3 with the above selection of $\xi_{m+1}$. Note that this strategy is closely related with [11], where the authors propose a greedy construction of a parameter-independent test space for Petrov-Galerkin projection, with a selection of basis functions based on an error indicator similar to $\delta_{r,m}(\xi)$.

4.3 Recycling factorizations of operator’s evaluations - Application to reduced basis method

When using a sample-based approach for solving a parameter-dependent equation (20), the linear system is solved for many values of the parameter $\xi$. When using a direct solver for solving a linear system for a given $\xi$, a factorization of the operator is usually available and can be used for improving a preconditioner for the solution of subsequent linear systems.

We here describe this idea in the particular context of greedy algorithms for Reduced Basis method, where the interpolation points $\xi_1, \ldots, \xi_r$ for the interpolation of the inverse $A(\xi)^{-1}$ are taken as the evaluation points $\xi_1^{RB}, \ldots, \xi_r^{RB}$ for the solution. At iteration $r$, having a preconditioner $P_r(\xi)$ and an approximation $u_r(\xi)$, a new interpolation point is defined such that

$$
\xi_{r+1}^{RB} \in \arg\max_{\xi \in \Xi} \|P_r(\xi)(A(\xi)u_r(\xi) - b(\xi))\|.
$$

Algorithm 4 describes this strategy.
Algorithm 4 Reduced Basis method with recycling of operator factorizations.

**Require:** $A(\xi), b(\xi), V,$ and $R$.

**Ensure:** Approximation $u_R(\xi)$.

1. Initialize $u_0(\xi) = 0, P_0(\xi) = I$;
2. for $r = 0$ to $R − 1$ do
3. Find $\xi_{r+1}^{RB} \in \arg\max_{\xi \in \Xi} \| P_r(\xi)(A(\xi)u_r(\xi) − b(\xi)) \|$
4. Compute a factorization of $A(\xi_{r+1}^{RB})$;
5. Solve the linear system $v_{r+1} = A(\xi_{r+1}^{RB})^{-1}b(\xi_{r+1}^{RB})$;
6. Update the approximation subspace $X_{r+1} = X_r + \text{span}(v_{r+1})$;
7. Define the implicit operator $P_{r+1} = A(\xi_{r+1}^{RB})^{-1}$;
8. Update the space $Y_{r+1}$ (or $Y_{r+1}^+$);
9. Compute the preconditioner: $P_{r+1}(\xi) = \arg\min_{P \in Y_{r+1} \text{ or } Y_{r+1}^+} \| (I - PA(\xi))V \|_F$;
10. Compute the Petrov-Galerkin approximation $u_{r+1}(\xi)$ of $u(\xi)$ on $X_{r+1}$ using equation (23);
11. end for

## 5 Numerical results

### 5.1 Illustration on a one parameter dependent model

In this section we compare the different interpolation methods on the following one parameter dependent advection-diffusion-reaction equation:

$$-\Delta u + v(\xi) \cdot \nabla u + u = f$$

defined over a square domain $\Omega = [0, 1]^2$ with periodic boundary conditions. The advection vector field $v(\xi)$ is spatially constant and depends on the parameter $\xi$ that takes values in $[0, 1]: v(\xi) = D \cos(2\pi \xi)e_1 + D \sin(2\pi \xi)e_2$, with $D = 50$ and $(e_1, e_2)$ the canonical basis of $\mathbb{R}^2$. $\Xi$ denotes a uniform grid of 250 points on $[0, 1]$. The source term $f$ is represented Figure 2(a). We introduce a finite element approximation space of dimension $n = 1600$ with piecewise linear approximations on a regular mesh of $\Omega$. A Galerkin projection on this approximation space yields the linear system of equations $A(\xi)u(\xi) = b$, with

$$A(\xi) = A_0 + \cos(2\pi \xi)A_1 + \sin(2\pi \xi)A_2,$$

where the matrices $A_0, A_1, A_2$ and the vector $b$ are given by

$$(A_0)_{i,j} = \int_{\Omega} \nabla \phi_i \cdot \nabla \phi_j + \phi_i \phi_j, \quad (A_1)_{i,j} = \int_{\Omega} \phi_i (e_1 \cdot \nabla \phi_j)$$

$$(A_2)_{i,j} = \int_{\Omega} \phi_i (e_2 \cdot \nabla \phi_j), \quad (b)_i = \int_{\Omega} \phi_i f,$$
where \( \{ \phi_i \}_{i=1}^n \) is the basis of the finite element space. Figures 2(b), 2(c) and 2(d) show three samples of the solution.

![Images of plots](a) f (b) u(0.05) (c) u(0.2) (d) u(0.8)

Figure 2: Plot of the source term \( f \) (a) and 3 samples of the solution corresponding to parameter values \( \xi = 0.05 \) (b), \( \xi = 0.2 \) (c) and \( \xi = 0.8 \) (d) respectively.

5.1.1 Comparison of the interpolating strategies

We first choose arbitrarily 3 interpolation points \( (\xi_1 = 0.05, \xi_2 = 0.2 \text{ and } \xi_3 = 0.8) \) and show the benefits of the use of the Frobenius norm projection for the definition of the preconditioner. For the comparison, we consider the Shepard and the nearest neighbor interpolation strategies. Let \( \| \cdot \|_{\Xi} \) denotes a norm on the parameter set \( \Xi \). The Shepard interpolation method is an inverse weighted distance interpolation:

\[
\lambda_i(\xi) = \begin{cases} 
\frac{\| \xi - \xi_i \|_{\Xi}^{-s}}{\sum_{j=1}^{m} \| \xi - \xi_j \|_{\Xi}^{-s}} & \text{if } \xi \neq \xi_i \\
1 & \text{if } \xi = \xi_i
\end{cases}
\]

where \( s > 0 \) is a parameter. Here we take \( s = 2 \). The nearest neighbor interpolation method consists in choosing the value taken by the nearest interpolation point, that means \( \lambda_i(\xi) = 1 \) for some \( i \in \arg \min_j \| \xi - \xi_j \|_{\Xi} \), and \( \lambda_j(\xi) = 0 \) for all \( j \neq i \).

Concerning the Frobenius norm projection on \( Y_m \) (or \( Y_m^+ \)), we first construct the affine decomposition of \( M(\xi) \) and \( S(\xi) \) as explained in Section 2.4. The interpolation points \( \xi_k^* \) (resp. \( \tilde{\xi}_k^* \)) given by the EIM procedure for \( M(\xi) \) (resp. \( S(\xi) \)) are \{0.0; 0.25; 0.37; 0.56; 0.80\} (resp. \{0.0; 0.25; 0.62\}). The number of terms \( m_M = 5 \) in the resulting affine decomposition of \( M(\xi) \) (see equation (19)) is less than the expected number \( m_A^2 = 9 \) (see equation (18a)). Considering the functions \( \Phi_1(\xi) = 1, \Phi_2(\xi) = \cos(2\pi \xi), \Phi_3(\xi) = \sin(2\pi \xi) \), and thanks to relation \( \cos^2 = 1 - \sin^2 \), the space

\[
\text{span}_{i,j} \{ \Phi_i \Phi_j \} = \text{span}\{1, \cos, \sin, \cos \sin, \cos^2, \sin^2\} = \text{span}\{1, \cos, \sin, \cos \sin, \cos^2\}
\]
is of dimension $m_M = 5$. The EIM procedure automatically detects the redundancy in the set of functions and reduces the number of terms in the decomposition \cite{19}. Then, since the dimension $n$ of the discretization space is reasonable, we compute the matrices $M(\xi_k^*)$ and the vectors $S(\xi_k^*)$ using equation \cite{19}.

The functions $\lambda_i(\xi)$ are plotted on Figure 3 for the proposed interpolation strategies. It is important to note that contrary to the Shepard or the nearest neighbor method, the Frobenius norm projection (on $Y_m$ or $Y_m^+$) leads to periodic interpolation functions, i.e. $\lambda_i(\xi = 0) = \lambda_i(\xi = 1)$. This is consistent with the fact that the application $\xi \mapsto A(\xi)$ is $1$-periodic. The Frobenius norm projection automatically detects such a feature.

![Figure 3: Interpolation functions $\lambda_i(\xi)$ for different interpolation methods.](image)

Figure 3 shows the condition number $\kappa_m(\xi)$ of $P_m(\xi)A(\xi)$ with respect to $\xi$. We first note that for the constant preconditioner $P_1(\xi) = A(\xi_2)^{-1}$, the resulting condition number is higher than the one of the non preconditioned matrix $A(\xi)$ for $\xi \in [0.55; 0.95]$. We also note that the Frobenius norm projection based interpolation strategies lead to better preconditioners than with the Shepard and nearest neighbor strategies. When considering the projection on $Y_m^+$ and $Y_m^\tilde{\kappa}$ (with $\tilde{\kappa} = 5 \times 10^4$ such that \cite{15} holds), the resulting condition number is roughly the same, so as the interpolation functions of Figures 3(d) and 3(e).

Since the projection on $Y_m^\tilde{\kappa}$ requires the expensive computation of the constants $\gamma^+, \gamma^-$ and $C$ (see Section 2.3), we prefer to simply use the projection on $Y_m^+$ in order to ensure the preconditioner to be invertible. Finally, for this example, it is not necessary to impose any
constraint since the projection on $Y_m$ leads to the best preconditioner and this preconditioner appears to be invertible for any $\xi \in \Xi$.

![Figure 4: Condition number of $P_m(\xi)A(\xi)$ for different interpolation strategies. The condition number of $A(\xi)$ is given as a reference.](image)

### 5.1.2 Using the Frobenius semi-norm

We analyze now the interpolation method defined by the Frobenius semi-norm projection on $Y_m$ for the different definitions of $V \in \mathbb{R}^{n \times K}$ proposed in sections 2.2.1 and 2.2.2. According to Table 1, the error on the interpolation functions decreases slowly with $K$ (roughly as $O(K^{-1/2})$), and the use of the P-SRHT matrix leads to a slightly lower error. The interpolation functions are plotted on Figure 5(a) in the case where $K = 8$. Even if we have an error of 36% to 101% on the interpolation functions, the condition number given on Figure 5(b) remains close to the one computed with the Frobenius norm. Also, an important remark is that with $K = 8$ the computational effort for computing $MV(\xi^*_k)$ and $SV(\tilde{\xi}^*_k)$ is negligible compared to the one for $M(\xi^*_k)$ and $S(\tilde{\xi}^*_k)$.

### 5.1.3 Greedy selection of the interpolation points

We now consider the greedy selection of the interpolation points presented in Section 4. We start with an initial point $\xi_1 = 0$ and the next points are defined by (31), where matrix $V$ is a realization of the P-SRHT matrix with $K = 128$ columns. $P_m(\xi)$ is the projection on $Y_m$ using the Frobenius semi-norm defined by (5). The first 3 steps of the algorithm are illustrated on Figure 6. We observe that at each iteration, the new interpolation point $\xi_{m+1}$
Table 1: Relative error $\sup_{\xi} \| \lambda(\xi) - \lambda^V(\xi) \|_{\mathbb{R}^2} / \sup_{\xi} \| \lambda(\xi) \|_{\mathbb{R}^2}$: $\lambda^V(\xi)$ (resp. $\lambda(\xi)$) are the interpolation functions associated to the Frobenius semi-norm projection (resp. the Frobenius norm projection) on $Y_m$, with $V$ either the rescaled partial Hadamard matrix, the random rescaled Rademacher matrix or the P-SRHT matrix (3 different samples for random matrices).

is close to the point where the condition number of $P_m(\xi)A(\xi)$ is maximal. Table 2 presents the maximal value over $\xi \in \Xi$ of the residual, and of the condition number of $P_m(\xi)A(\xi)$. Both quantities are rapidly decreasing with $m$. This shows that this algorithm, initially designed to minimize $\|(I - P_m(\xi)A(\xi))V\|_F$, seems to be also efficient for the construction of preconditioners, in the sense that the condition number decreases rapidly.

| $K$             | 8   | 16  | 32  | 64  | 128 | 256 | 512 |
|-----------------|-----|-----|-----|-----|-----|-----|-----|
| Rescaled partial Hadamard | 0.4131 | 0.3918 | 0.3221 | 0.1010 | 0.0573 | 0.0181 | 0.0255 |
| Rescaled Rademacher (1)     | 0.5518 | 0.0973 | 0.2031 | 0.1046 | 0.1224 | 0.1111 | 0.0596 |
| Rescaled Rademacher (2)     | 1.0120 | 0.6480 | 0.1683 | 0.1239 | 0.0597 | 0.0989 | 0.0514 |
| Rescaled Rademacher (3)     | 0.7193 | 0.2014 | 0.1241 | 0.1051 | 0.1235 | 0.1369 | 0.0519 |
| P-SRHT (1)           | 0.4343 | 0.2081 | 0.2297 | 0.0741 | 0.0723 | 0.0669 | 0.0114 |
| P-SRHT (2)           | 0.3624 | 0.2753 | 0.0931 | 0.1285 | 0.0622 | 0.0649 | 0.0249 |
| P-SRHT (3)           | 0.8133 | 0.4227 | 0.1138 | 0.0741 | 0.0824 | 0.0469 | 0.0197 |

Table 2: Convergence of the greedy algorithm: supremum over $\xi \in \Xi$ of the condition number (first row) and of the Frobenius semi-norm residual (second row).

5.2 Multi-parameter dependent equation

We introduce a benchmark proposed within the OPUS project (see http://www.opus-project.fr). Two electronic components $\Omega_{IC}$ (see Figure 7) submitted to a cooling air flow in the domain $\Omega_{Air}$ are fixed on a printed circuit board $\Omega_{PCB}$. The temperature field defined over $\Omega = \Omega_{IC} \cup \Omega_{PBC} \cup \Omega_{Air} \subset \mathbb{R}^2$ satisfies the advection-diffusion equation:

$$-\nabla \cdot (\kappa(\xi) \nabla u) + D(\xi)v \cdot \nabla u = f. \quad (35)$$

The diffusion coefficient $\kappa(\xi)$ is equal to $\kappa_{PCB}$ on $\Omega_{PCB}$, $\kappa_{Air}$ on $\Omega_{PCB}$ and $\kappa_{IC}$ on $\Omega_{IC}$. The right hand side $f$ is equal to $Q = 10^6$ on $\Omega_{IC}$ and 0 elsewhere. The boundary conditions are $u = 0$ on $\Gamma_d$, $e_2 \cdot \nabla u = 0$ on $\Gamma_u$ ($e_1, e_2$ are the canonical vectors of $\mathbb{R}^2$), and $u|_{\Gamma_l} = u|_{\Gamma_r}$.
Figure 5: Comparison between the Frobenius norm projection on $Y_3$ (black lines) and the Frobenius semi-norm projection on $Y_3$, using for $V$ either a sample of the rescaled Rademacher matrix (blue lines), the rescaled partial Hadamard matrix (red lines) or a sample of the P-SRHT matrix (green lines) with $K = 8$.

Figure 6: Greedy selection of the interpolation points: the first row is the residual $\| (I - P_k(\xi)A(\xi))V \|_F$ (the blue points correspond to the maximum of the residual) with $V$ a realization of the P-SRHT matrix with $K = 128$ columns, and the second row is the condition number of $P_m(\xi)A(\xi)$.
(periodic boundary condition). At the interface \( \Gamma_C = \partial \Omega_{IC} \cap \partial \Omega_{PCB} \) there is a thermal contact conductance, meaning that the temperature field \( u \) admits a jump over \( \Gamma_C \) which satisfies

\[
\kappa_{IC}(e_1 \cdot \nabla u|_{\Omega_{IC}}) = \kappa_{PCB}(e_1 \cdot \nabla u|_{\Omega_{PCB}}) = r(u|_{\Omega_{IC}} - u|_{\Omega_{PCB}}) \quad \text{on } \Gamma_C.
\]

The advection field \( v \) is given by \( v(x,y) = e_2 g(x) \), where \( g(x) = 0 \) if \( x \leq e_{PCB} + e_{IC} \) and

\[
g(x) = \frac{3}{2(e - e_{IC})} \left(1 - \left(\frac{2x - (e + e_{IC} + 2e_{PCB})}{e - e_{IC}}\right)^2\right)
\]

otherwise. We have 4 parameters: the width \( e := \xi_1 \) of the domain \( \Omega_{Air} \), the thermal conductance parameter \( r := \xi_2 \), the diffusion coefficient \( \kappa_{IC} := \xi_3 \) of the components and the amplitude of the advection field \( D := \xi_4 \). Since the domain \( \Omega = \Omega(e) \) depends on the parameter \( \xi_1 \in [e_{min}, e_{max}] \), we introduce a geometric transformation \( (x,y) = \phi_{\xi_1}(x_0,y_0) \) that maps a reference domain \( \Omega_0 = \Omega(e_{max}) \) to \( \Omega(\xi_1) \):

\[
\phi_{\xi_1}(x_0,y_0) = \begin{cases} x_0 & \text{if } x_0 \leq e_0 \\ e_0 + (x_0 - e_0) \frac{\xi_1 - e_{IC}}{e_{max} - e_{IC}} & \text{otherwise.} \end{cases}
\]

with \( e_0 = e_{PCB} + e_{IC} \). This method is described in [29]: since the geometric transformation \( \phi_{\xi_1} \) satisfies the so called Affine Geometry Precondition, the operator of equation (35) formulated on the reference domain admits an affine representation.

For the spatial discretization we use a finite element approximation with \( n = 2.8 \cdot 10^4 \) degrees of freedom (piecewise linear approximation). We rely on a Galerkin method with SUPG stabilization (see [7]). \( \Xi \) is a set of \( 10^4 \) independent samples drawn according the lognormal probability laws of the parameters given on Figure 7.

### 5.2.1 Preconditioner for the projection on a given reduced space

We consider here a POD basis \( X_r \) of dimension \( r = 50 \) computed with 100 snapshots of the solution (a basis of \( X_r \) is obtained by the first 50 dominant singular vectors of a matrix of 100 random snapshots of \( u(\xi) \)). Then we compute the Petrov-Galerkin projection as presented in Section 3.1. The efficiency of the preconditioner can be measured with the quantity \( \delta_{r,m}(\xi) \): the associated quasi-optimality constant \((1 - \delta_{r,m}(\xi)^2)^{-1/2} \) should be as close to one as possible (see equation (26)). We introduce the quantile \( q_p \) of probability \( p \) associated to the quasi-optimality constant \((1 - \delta_{r,m}(\xi)^2)^{1/2} \) defined as the smallest value \( q_p \geq 1 \) satisfying

\[
P\left\{\{\xi \in \Xi : (1 - \delta_{r,m}(\xi)^2)^{1/2} \leq q_p\}\right\} \geq p,
\]

28
where \( P(A) = \#A/\#\Xi \) for \( A \subseteq \Xi \). Table 3 shows the evolution of the quantile with respect to the number of interpolation points for the preconditioner. Here the goal is to compare the different strategies for the selection of the interpolation points:

(a) the greedy selection (34) based on the quantity \( \delta_{r,m}(\xi) \),

(b) the greedy selection (31) based on the Frobenius semi-norm residual, with \( V \) the rescaled partial Hadamard matrix with \( K = 256 \) columns, and

(c) a random Latin Hypercube sample (LHS).

The projection on \( Y_m \) (or \( Y_m^+ \)) is then defined with the Frobenius semi-norm using for \( V \) a P-SRHT matrix with \( K = 330 \) columns.

When considering a small number of interpolation points \( m \leq 3 \), the projection on \( Y_m^+ \) provides lower quantiles for the quasi-optimality constant compared to the projection on \( Y_m \). The positivity constraint is useful for small \( m \). But for high values of \( m \) (see \( m = 15 \)) the positivity constraint is no longer necessary and the projection on \( Y_m \) provides lower quantiles.

Concerning the choice of the interpolation points, the strategy (a) shows the faster decay of the quantiles \( q_p \), especially for \( p = 50\% \). The strategy (b) shows also good results, but the quantile \( q_p \) for \( p = 100\% \) are still high compared to (a). These results show the benefits of the greedy selection based on the quasi-optimality constant. Finally the strategy (c) shows bad results (high values of the quantiles), especially for small \( m \).

### 5.2.2 Preconditioner for Reduced Basis method

We now consider the preconditioned Reduced Basis method for the construction of the approximation space \( X_r \), as presented in Section 3.2. Figures 9 and 10 show the convergence

---

**Figure 7**: Geometry and parameters of the benchmark OPUS.
Table 3: Quantiles \( q_p \) of the quasi-optimality constant associated to the Petrov-Galerkin projection on the POD subspace \( X_r \) for \( p = 50\% \), 90\% and 90\%. The row \( m = 0 \) corresponds to \( P_0(\xi) = R_X^{-1} \), that is the standard Galerkin projection.

of the error with respect to the rank \( r \) of \( u_r(\xi) \) for different constructions of the preconditioner \( P_m(\xi) \). Two measures of the error are given: \( \sup_{\xi \in \Xi} \| u(\xi) - u_r(\xi) \| / \| u(\xi) \| \), and the quantile of probability 0.97 for \( \| u(\xi) - u_r(\xi) \| / \| u(\xi) \| \). The curve “Ideal greedy” corresponds to the algorithm defined by [27] which provides a reference for the ideally conditioned algorithm, i.e. with \( \kappa_m(\xi) = 1 \). Figure 8 shows the corresponding first interpolation points for the solution.

The greedy selection of the interpolation points based on [31] (see Figure 9) allows to almost recover the convergence curve of the ideal greedy algorithm when using the projection on \( Y_m \) with \( m = 15 \). For the recycling strategy, the approximation is rather bad for \( r = m \leq 10 \) meaning that the space \( Y_r \) (or \( Y_r^+ \)) is not really adapted for the construction of a good preconditioner over the whole parametric domain. However, for higher values of \( r \), the preconditioner is getting better and better. For \( r \geq 20 \), we almost reach the convergence of the ideal greedy algorithm. We conclude that this recycling strategy, with
a computational cost which is comparable to the standard non preconditioned Reduced Basis greedy algorithm, allows obtaining asymptotically the performance of the ideal greedy algorithm. Note that the positivity constraint yields a better preconditioner for small values of $r$ but is no longer necessary for large $r$.

Let us finally consider the effectivity index

$$\eta_r(\xi) = \frac{\|P_r(\xi)(A(\xi)u_r(\xi) - b(\xi))\|_{X'}}{\|u(\xi) - u_r(\xi)\|_X},$$

which evaluates the quality of the preconditioned residual norm for error estimation. We introduce the confidence interval $I_r(p)$ defined as the smallest interval which satisfies

$$\mathbb{P}(\{\xi \in \Xi : \eta_r(\xi) \in I_r(p)\}) \geq p.$$

On Figure 8, we see that the confidence intervals are shrinking around 1 when $r$ increases, meaning that the preconditioned residual norm becomes a better and better error estimator when $r$ increases. Again, the positivity constraint is needed for small values of $r$, but we obtain a better error estimation without imposing this constraint for $r \geq 20$. On the contrary, the standard residual norm leads to effectivity indices that spread from $10^{-1}$ to $10^1$ with no improvement as $r$ increases, meaning that we can have a factor $10^2$ between the error estimator $\|A(\xi)u_r(\xi) - b(\xi)\|_{X'}$ and the true error $\|u_r(\xi) - u(\xi)\|_X$. 

Let us finally consider the effectivity index

$$\eta_r(\xi) = \frac{\|P_r(\xi)(A(\xi)u_r(\xi) - b(\xi))\|_{X'}}{\|u(\xi) - u_r(\xi)\|_X},$$

which evaluates the quality of the preconditioned residual norm for error estimation. We introduce the confidence interval $I_r(p)$ defined as the smallest interval which satisfies

$$\mathbb{P}(\{\xi \in \Xi : \eta_r(\xi) \in I_r(p)\}) \geq p.$$

On Figure 8, we see that the confidence intervals are shrinking around 1 when $r$ increases, meaning that the preconditioned residual norm becomes a better and better error estimator when $r$ increases. Again, the positivity constraint is needed for small values of $r$, but we obtain a better error estimation without imposing this constraint for $r \geq 20$. On the contrary, the standard residual norm leads to effectivity indices that spread from $10^{-1}$ to $10^1$ with no improvement as $r$ increases, meaning that we can have a factor $10^2$ between the error estimator $\|A(\xi)u_r(\xi) - b(\xi)\|_{X'}$ and the true error $\|u_r(\xi) - u(\xi)\|_X$. 

Figure 8: First six interpolation points of the ideal reduced basis method and corresponding reduced basis functions.
Figure 9: Convergence of the preconditioned reduced basis method using the greedy selection of interpolation points for the preconditioner. Supremum over Ξ (top) and quantile of probability 97% (bottom) of the relative error $\|u(\xi) - u_r(\xi)\|_X / \|u(\xi)\|_X$ with respect to $r$. Comparison of preconditioned reduced basis algorithms with ideal and standard greedy algorithms.

6 Conclusion

We have proposed a method for the interpolation of the inverse of a parameter dependent matrix. The interpolation is defined by the projection of the identity in the sense of the
Figure 10: Preconditioned Reduced basis methods with recycling. Supremum over $\Xi$ (left) and quantile of probability 97% (right) of the relative error $\|u(\xi) - u_r(\xi)\|_X / \|u(\xi)\|_X$ with respect to $r$. Comparison of preconditioned reduced basis algorithms with ideal and standard greedy algorithms.

Figure 11: Confidence intervals of the effectivity index during the iterations of the Reduced Basis greedy construction. Comparison between preconditioned algorithms with recycling of operators factorizations (a,b) and the non preconditioned greedy algorithm (c).

Frobenius norm. Approximations of the Frobenius norm (based on the use of Hadamard matrices or random matrices) have been introduced to make computationally feasible the projection in the case of large matrices. Different strategies have then been proposed for the selection of interpolation points depending of the objective: (i) the construction of an
optimal approximation of the inverse operator for preconditioning iterative solvers or for improving error estimators based on preconditioned residuals, (ii) the improvement of the quality of Petrov-Galerkin projections of the solution of a parameter-dependent equation on a given reduced approximation space, or (iii) the recycling of operator factorizations when solving a parameter-dependent equation with a sample-based approach. The performance of the obtained parameter-dependent preconditioner has been illustrated in the context of projection based model reduction techniques such as the Proper Orthogonal Decomposition and the Reduced Basis method.

In this paper, we have restricted the presentation to the case of real matrices but the methodology can be naturally extended to the case of complex matrices.

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