Solving Optimal Power Flow Using Modified Bacterial Foraging Algorithm Considering FACTS Devices

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Abstract

In this paper, a new Modified Bacterial Foraging Algorithm (MBFA) method is developed to incorporate FACTS devices in optimal power flow (OPF) problem. This method can provide an enhanced economic solution with the use of controllable FACTS devices. Two types of FACTS devices, thyristor controlled series compensators (TCSC) and Static VAR Compensator (SVC) are considered in this method. The basic bacterial foraging algorithm (BFA) is an evolutionary optimization technique inspired by the foraging behavior of the E. coli bacteria. The strategy of the OPF problem is decomposed in two sub-problems, the first sub-problem related to active power planning to minimize the fuel cost function, and the second sub-problem designed to make corrections to the voltage deviation and reactive power violation based in an efficient reactive power planning of multi Static VAR Compensator (SVC). The specified power flow control constraints due to the use of FACTS devices are included in the OPF problem. The proposed method decomposes the solution of such modified OPF problem into two sub problems’ iteration. The first sub problem is a power flow control problem and the second sub problem is a modified Bacterial foraging algorithm (MBFA) OPF problem. The two sub problems are solved iteratively until convergence. Case studies are presented to show the effectiveness of the proposed method.

Keywords

Flexible AC Transmission System (FACTS); Modified Bacterial Foraging Algorithm (MBFA); Optimal Power Flow (OPF); TCSC; SVC

1. Introduction

The solution of an OPF problem determines the optimal settings for control variables in a power network observing various constraints. In the last 30 years’ development, many different solution approaches have been proposed to solve the OPF problems [1]. Recent developments in power electronics have opened a new world in power systems control. Several control devices are being developed under this new concept, known as FACTS
These FACTS devices can make the power systems operating in a more flexible, more secure and economic way. However, these FACTS devices also make the power system operating in a more complicated way. In [3], the Benders decomposition method was first proposed to solve the optimal active power flow dispatch problem incorporating FACTS devices. This method can deal with the representation of series compensators and phase shifters, but this method did not consider the specified line flow constraints. In [4], a method for solving the power flow control problem incorporating FACTS devices was proposed based on decomposition. However, this method did not combine tilt OPF problem with the power flow control problem, hence the solution probably may not be the overall optimal solution. In [5] author presented the solution of optimal power flow using linear programming method for power system network security application. Gotham and Heydt [6] have presented the modeling of FACTS devices for power flow studies and the role of that modeling in the study of FACTS devices. The presence of discrete control variables, such as switchable shunt devices, transformer tap positions, and phase shifters, further complicates the problem solution [7,8]. The global optimization techniques known as genetic algorithms (GA) [9], simulated annealing (SA), tabu search (TS) [10], evolutionary programming (EP), improved particle swarm optimization (IPSO) [11], a novel crazy swarm optimization (NCSO) [12] and a hybrid genetic algorithm approach based on differential evolution (IGA-DE) [6], in [13] the authors proposed a simple genetic algorithm fuzzy system and evolutionary programming applied to the OPF problem in large-scale power systems. Author in [14] presents a particle swarm optimization (PSO) to solve the economic dispatch with consideration of many nonlinear characteristics of the generator, such as ramp rate limits, prohibited operating zones, and non-smooth cost functions. In [15] authors present a novel string structure for solving the economic dispatch through genetic algorithm (GA) with consideration of practical generators constraints. In [16] authors proposed optimal power dispatch for large scale power system using stochastic search algorithms. To accelerate the search process authors in [10] proposed a multiple tabu search algorithm. To overcome the drawbacks of the conventional methods related to the form of the cost function, and to reduce the computational time related to the large space search required by GA, this paper presents a modified BFA for the solution of large-scale OPF with practical generators constraints and with consideration shunt FACTS devices.

This paper is organized as follows: Modeling of FACTS device is given in Section 2; Problem Formulation is given in Section 3. Modified Bacterial Foraging algorithm for proposed method is given in Section 4. Results and discussion are given in Section 5, the conclusion is drawn in Section 6 and references are given in Section 7.

2. Modeling of Fact Device

2.1. Thyristor Controlled Series Compensator (TCSC)

When the DC network model is used, the model of the transmission line with TCSC is shown in $s$. The total susceptance of the transmission line can be formulated as:

$$\gamma_s = \frac{1}{2}(x_g - x_i)$$  \hspace{1cm} (1)

Which is directly used as the control variable to be implemented in the bus susceptance matrices (Figure 1).

The power flow equations of the branch can be formulated as follows:

$$P_{ij} = \gamma_s \delta_{ij}$$  \hspace{1cm} (2)

2.2. Static VAR Compensator (SVC)

The steady-state proposed model is used to incorporate the SVC on power flow problems. This model is based on representing the controller as a variable impedance, assuming an SVC configuration with a fixed capacitor (FC) and Thyristor-controlled reactor (TCR) as depicted in Figure 2. Applying simultaneously a gate pulse to

![Figure 1. Equivalent circuit of TCSC.](image-url)
all thyristor of a thyristor valve brings the valve into conduction. The valve will block approximately at the zero crossing of the ac current, in the absence of firing signals. Thus, the controlling element is the Thyristor valve. The thyristors are fired symmetrically, in an angle control range of 90 - 180 with respect to the capacitor (inductor) voltage.

\[ V = V_{ref} + X_{sl} \]  \hspace{1cm} (3)

\( X_{sl} \) are in the range of 0.02 - 0.05 p.u. with respect to the SVC base. The slope is needed to avoid hitting limits. At the voltage limits the SVC is transformed into a fixed reactance. The total equivalent impedance \( X_e \) of SVC may be represented by

\[ X_e = X_C \frac{\pi / k_s}{\sin 2\alpha - 2\alpha + \pi (2 - 1/k_s)} \]  \hspace{1cm} (4)

where \( k_s = X_C / X_L \).

3. Problem Formulation

In this paper, as our attention is focused on the optimal active power flow problem, the liberalized (DC) power flow model is used. When the \( n - 1 \) security constraints were considered, the OPF problem can be formulated as follow:

To minimize \( CP \),

\[ \sum P_i = \sum L_i + P_i \]  \hspace{1cm} (5)

\[ SF + P = L \]

\[ F - \gamma S' \delta = 0 \]  \hspace{1cm} (6)

\[ P \leq P \leq \bar{P} \]

\[ F \leq F(0) \leq F \]

\[ F \leq F(l) \leq \bar{F} \]

where,

- \( C \): is the generation cost vector
- \( P \): is the vector of active bus generation
- \( s' \): is the transpose of bus-branch incidence matrix
- \( \gamma \): is the diagonal matrix of circuit susceptance
- \( \delta \): is the vector of angle across branch
- \( F(0) \) is the vector of base state active power flow in lines
- \( F(l) \) is the vector of base state active power flow in lines when branch \( l \) is break out
- \( P \): is the active power loss in the system
- \( L \): is the vector of active bus loads
- \( F \) and \( \bar{F} \) is the circuit flow limits
- \( P \) and \( \bar{P} \) are the generation limits

This conventional optimization problem can be solved by the LP-based algorithm efficiently [5,17]. Next, two additional control variables and a set of power flow control equality constraints will be introduced into this optimization.
Because the FACTS devices parameters can be controlled in very small steps, the new additional control variables can be seen as continuous and limited by lower and upper bounds. The modified optimal active power flow problem can be formulated as follows:

\[ \sum P = \sum L_i + P_i \]  \hspace{1cm} (7)
\[ SF + P = L \]  \hspace{1cm} (8)
\[ F - \gamma S' \delta = 0 \]
\[ W^\nu = \gamma c(s, \delta + \Psi) \]  \hspace{1cm} (9)
\[ P \leq P_c \leq \bar{P} \]
\[ F \leq F(0) \leq \bar{F} \]
\[ F \leq F(l) \leq \bar{F} \]
\[ \gamma \leq \gamma \leq \bar{\gamma} \]
\[ \Psi \leq \Psi \leq \Psi \]

where, \( \bar{\gamma} \) is the diagonal matrix of controlled line susceptance and \( SC \) is the bus-branch incidence matrix related to controlled line. While \( \Psi \) is the vector of controlled line phase shifter angle and \( W^\nu \) is a vector of specified line flows. \( P_c \) is the vector of bus generation including the compensation injection power of phase shifter. It can be seen that this OPF problem is no longer a linear optimization problem due to the new control variable \( \gamma \) and the specified power flow equality constraints. Thus, it is not possible to use conventional LP technique [5,17] directly, and it becomes necessary to develop a new method to solve this optimization problem.

4. The Proposed Mbfa

4.1. Bacterial Foraging Algorithm

BFA is an optimization technique motivated by the foraging behavior of the \textit{E coli}. Bacteria. The biological aspects of the bacterial foraging strategies and their motile behavior as well as their decision making mechanisms can be found in [9]. BFA is designed to solve non-gradient optimization problems and to handle complex and non-differentiable objective functions. Searching the hyperspace is performed through three main operations, namely; chemotaxis, reproduction and elimination dispersal activities [18]. The chemotaxis process is performed through swimming and tumbling. The bacterium spends its life alternating between these two modes of motion. In the BFA, a tumble is represented by a unit length in a random direction, \( \Phi(j) \), which specifies the direction of movement after a tumble. The size of the step taken in the random direction is represented by the constant run-length unit, \( C(i) \). For a population of bacteria, the location of the \( i^{th} \) bacterium at the \( j^{th} \) chemotactic step, \( k^{th} \) reproduction step and \( l^{th} \) elimination/dispersal event is represented by \( \theta(i,j,k,l) \in \mathbb{R}^n \).

At this location the cost function is denoted by

\[ \theta'(j + 1,k,l) = \theta'(j,k,l) + C(i,j)\phi(j) \]  \hspace{1cm} (10)

When, \( \theta'(j + 1,k,l) \) the cost function \( J(i, j + 1, k, l) \) better (lower) than \( J(i,j,k,l) \), another step of size \( C(i,j) \) in the same direction is taken. This swimming operation is repeated as long as a lower cost is obtained until a maximum preset number of steps, \( N_s \), is reached. The cost function of each bacterium in the population is affected by a kind of swarming that is performed by the cell-to-cell signaling released by the bacteria groups to form swarm patterns. This swarming is expressed as follows:

\[ J_s(\theta, P(j,k,l)) = \sum_{j=1}^{p} \sum_{k=1}^{m} \sum_{l=1}^{n} h_{\text{attract}} \sum_{m=1}^{p} (\theta_m - \theta_m')^2 \]

\[ = \sum_{j=1}^{p} \sum_{k=1}^{m} \sum_{l=1}^{n} h_{\text{repellant}} \sum_{m=1}^{p} (\theta_m - \theta_m')^2 \]

where \( h_{\text{attract}}, \omega_{\text{attract}}, h_{\text{repellant}} \) and \( \omega_{\text{repellant}} \) are coefficients that represent the characteristics of the
attractant and repellant signals released by the cell and \( \theta_m \) is the \( m \)th component of \( i \)th bacterium position \( P(i,j,k,l) \) is the position of each member of the population of the S bacteria and defined as:

\[
P(j,k,l) = \{\theta^t(j,k,l) | i = 1,2,...,S\}
\]

where \( S \) is the size of the bacteria population. The function (9) which represents the cell-to-cell signaling effect is added to the cost function

\[
J(i,j,k,l) + J_c(\theta, P)
\]

A reproduction process is performed after taking a maximum number of chemotactic steps, \( N_c \). The population is halved so that the least healthy half dies and each bacterium in the other healthiest one splits into two bacteria which takes the same position.

\[
S_2 = \frac{S}{2}
\]

After \( N_{re} \) reproduction steps an elimination/dispersal event takes place for \( N_{ed} \) number of excisions. In this operation each bacterium could be moved to explore other parts of the search space. The probability for each bacterium to experience the elimination/dispersal event is determined by a predefined fraction \( P_{ed} \).

4.2. Modified Bacterial Foraging Algorithm

The unit step length of the basic BFA is constant which may guarantee good searching results for all optimization problems. However, when applied to complex problems with high dimensionality it shows poor performance. The run length parameter is the key factor for controlling the local and global search ability of the BFA. From this perspective, balancing the exploration and exploitation of the search could be achieved by adjusting the run-length unit. In this paper we propose a nonlinear decreasing dynamic function to perform the swim walk instead of the constant step. This function is expressed as:

\[
C(i,j+1) = \left(\frac{C(i,j) - C(N_c)}{N_c + C(N_c)}\right)(N_c - j)
\]

where \( j \) is the chemotactic step and \( N_c \) is the maximum number of chemotactic steps while \( C(N_c) \) is a predefined parameter. Control parameters of the MBFA are given in Table 1.

The algorithm of the proposed technique is as follows:

**Step 1:** Initialization of the following parameters:
- \( p \): Dimension of the search space
- \( S \): The number of bacteria in the population
- \( N_c \): Number of chemotactic steps
- \( N_s \): The length of a swim when it is on a gradient
- \( N_{re} \): The number of reproduction steps
- \( N_{ed} \): The number of elimination-dispense steps
- \( P_{ed} \): The probability that each bacterium will be eliminated/dispersed

\( C(i,j) = \): Initial run-length unit

\( C(N_c) \): The run-length unit at the end of the

| Sl. No | PARAMETER | VALUE |
|-------|-----------|-------|
| 1     | Number of bacteria’s | 4     |
| 2     | Maximum number of steps, \( N_c \) | 4     |
| 3     | Number of chemotactic steps, \( N_c \) | 5     |
| 4     | Number of reproduction steps, \( N_{re} \) | 4     |
| 5     | Number of elimination-dispense steps, \( N_{ed} \) | 2     |
| 6     | Probability, \( P_{ed} \) | 0.2   |
| 7     | Size of the step, \( C(i) \) | 0.1   |
chemotactic steps \( j = N_c \).

\( \theta^i \): The initial random location of each bacterium

Step 2: Elimination/dispersal loop, \( l = l + 1 \)

Step 3: Reproduction loop, \( k = k + 1 \)

Step 4: Chemotaxis loop, \( j = j + 1 \)

• For \( i = 1, 2, \ldots, S \), execute the chemotactic step for each bacterium as follows:

- Evaluate the cost function \( J(i, j, k, l) \) using (9) and (11).
- Let \( J_{\text{last}} = J(i, j, k, l) \) so that a lower cost could be found.
- Tumble: generate a random vector \( \Delta(i) \in \mathbb{R}^p \) and \( \Delta_m(i), m = 1, 2, \ldots, p \) is a random number in the range \([-1, 1]\).
- Compute \( \phi(i) = \frac{\Delta(i)}{\sqrt{\sum \Delta^2(i)\Delta(i)}} \)
- Move using (8)
- Compute \( J(i, j+1, k, l) \) and use (9) to Compute \( J_u(\theta, P(j+1, k, l)) \) then use (11) to find the new \( J(i, j+1, k, l) \)

- Swim: let \( m = 0 \) (counter for swim length)

While \( m < N_s \) (no climbing down too long)

Let \( m = m + 1 \).

If \( J(i, j+1, k, l) < J_{\text{last}} \) let \( J_{\text{last}} = J(i, j+1, k, l) \) then take another step in the same direction and compute the new \( J(i, j+1, k, l)_{\text{last}} \)

- Go to the next bacterium \( i = i + 1 \) if \( i \neq S \).

- Update the run length using (15).

Step 5: If \( j < N_c \) go to step 4 \( (j = j + 1) \).

Step 6: Reproduction:

• For the given \( k, l \), evaluate the health of each bacterium \( i \) as follows:

\[
J^{i}_{\text{health}} = \sum_{j=1}^{N_c+1} J(i, j, k, l) \tag{16}
\]

health of the bacterium \( i \) measures how many nutrient it got over its lifetime.

• Sort bacteria according to their health \( J^{i}_{\text{health}} \) in ascending order.

• The bacteria with the highest \( J^{i}_{\text{health}} \) values, computed by (16) die while the other \( S \), with the lowest values split and take the same location of their parents.

Step 7: if \( k < N_r \), go to step 3 \( (k = k + 1) \)

Step 8: Elimination/ dispersal

• With probability \( ped \), randomly eliminate and dispersal each bacterium \( i \), keeping the size of the population constant.

Step 9: if \( l < N \), go to step 2 \( (l = l + 1) \), otherwise end

5. Simulation Results

The proposed algorithm is developed in the Matlab programming language using 7.5 versions. The proposed MBFA is tested using modified IEEE 30-Bus system. The test example has been run on a 2.6-Ghz Pentium-IV PC.

Case Studies on the IEEE 30-Bus System

IEEE 30-Bus, 41-branch system is tested, for the voltage constraint the lower and upper limits are 0.9 p.u and 1.06 p.u., respectively, (expect for PV buses where \( V_{\text{max}} = 1:1 \) p.u).

The efficiency of the proposed approach, we made a comparison of our algorithm with others competing OPF algorithm. In [9], they presented a standard GA, in [19], the authors presented an enhanced GA, and then in [20], they proposed an Ant colony optimization (ACO).

In [13] they presented an optimal power flow solution using GA-fuzzy system approach. In [21], they presented an efficient parallel GA (EPGA). The operating cost in our proposed approach is 800.1585 and the power loss is 8.4625 which are better than other methods reported in the literature. Results in Table 2 show clearly that
Table 2. Results of minimum cost and power generation with SGA, EGA, ACO, FGA and EPGA for IEEE 30 bus.

| Variable | SGA [9]  | EGA [19] | ACO [20] | FGA [13] | EPGA [21] | Proposed method (MBFA) |
|----------|----------|----------|----------|----------|----------|------------------------|
| P1 (MW)  | 179.367  | 176.20   | 181.945  | 175.137  | 180.12   | 164.501                |
| P2 (MW)  | 44.24    | 48.75    | 47.0010  | 50.353   | 44.18    | 32.423                 |
| P5 (MW)  | 24.61    | 21.44    | 20.5530  | 21.451   | 19.64    | 21.459                 |
| P8 (MW)  | 19.90    | 21.95    | 21.1460  | 21.176   | 20.96    | 11.130                 |
| P11 (MW) | 10.71    | 12.42    | 10.4330  | 12.667   | 14.90    | 23.872                 |
| P13 (MW) | 14.09    | 12.02    | 12.1730  | 12.11    | 12.72    | 38.470                 |
| Q1 (MVAR)| −3.156   | -        | −6.562   | −4.50    | −3.609   |                        |
| Q2 (MVAR)| 42.543   | -        | 22.356   | 30.71    | 37.423   |                        |
| Q5 (MVAR)| 26.292   | -        | 30.372   | 22.59    | 24.985   |                        |
| Q8 (MVAR)| 22.768   | -        | 18.89    | 37.85    | 19.530   |                        |
| Q11 (MVAR)| 29.923  | -        | 21.737   | −2.52    | 15.821   |                        |
| Q13 (MVAR)| 32.346  | -        | 22.635   | −13.08   | 7.598    |                        |
| θ1 (deg)| 0.000    | -        | 0.00     | 0.00     | 0.00     |                        |
| θ2 (deg)| −3.674   | -        | −3.608   | −3.448   | −3.375   |                        |
| θ5 (deg)| −10.14   | -        | −10.509  | −9.858   | −9.596   |                        |
| θ8 (deg)| −10.00   | -        | −8.154   | −7.638   | −7.012   |                        |
| θ11 (deg)| −8.851  | -        | −8.783   | −7.507   | −4.789   |                        |
| θ13 (deg)| −10.13   | -        | −10.228  | −9.102   | −4.584   |                        |
| Ploss (MW)| 9.5177  | 9.3900   | 9.8520   | 9.494    | 9.120    | 8.4625                 |
| Cost ($/h)| 803.699 | 802.06   | 802.578  | 802.0003 | 801.3445 | 800.1585               |

the proposed approach gives better results and the best solution of shunt compensation obtained at the standard load demand (Pd = 283.4 MW) using reactive power planning [18,22], line flows obtained are well under security limits compared to other algorithm such as FGA, SGA, ACO and other algorithms.

First of all, the total real power loss of the system is calculated with shunt FACTS device and the result is 8.4625 MW. And to improve the performance of reactive power planning of the system shunt FACTS device included at different locations. The power loss is optimized by the MBFA with the combination of improving the bus voltage profile and generator reactive power.

6. Conclusion

A modified bacterial foraging algorithm (MBFA) for solving large scale OPF Problem has been demonstrated in this paper. Simulation results have demonstrated the effectiveness of the proposed algorithm and Comparison with other computational methods has shown that the MBFA achieved better results. A new method to incorporating FACTS devised in OPF problem has been presented in this paper. The proposed method separates the modified optimization problem into problems’ iteration. The power flow control problem while the sub problem is a conventional OPF problem. The method incorporates the power flow control needs due to FACTS devices into the OPF Problem, and the optimization results can satisfy the power flow control needs and the economic operation at the same time. The case studies of an adapted IEEE 30-bus system have been presented to show the efficiency of the proposed method.

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