Comparison of graph clustering methods for analyzing the mathematical subject classification codes

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Abstract

Various graph clustering methods have been introduced to identify communities in social or biological networks. This paper studies the entropy-based and the Markov chain-based methods in clustering the undirected graph. We examine the performance of two clustering methods with conventional methods based on quality measures of clustering. For the real applications, we collect the mathematical subject classification (MSC) codes of research papers from published mathematical databases and construct the weighted code-to-document matrix for applying graph clustering methods. We pursue to group MSC codes into the same cluster if the corresponding MSC codes appear in many papers simultaneously. We compare the MSC clustering results based on the several assessment measures and conclude that the Markov chain-based method is suitable for clustering the MSC codes.

Keywords: mathematical subject classification, Markov chain clustering, entropy graph clustering

1. Introduction

The mathematical subject classification (MSC) is the five-digit alphanumerical code for the organization of research papers, which are based on two mathematical databases, mathematical reviews (MR) and Zentralblatt MATH (zbMath) (Lange \textit{et al.}, 2012). The primary purpose of the use of MSC is to help users find the literature of interest by subject area (https://msc2020.org/). Various mathematical and other scientific journals often request authors to enter the MSC codes in their articles, and the author-supplied MSC codes are used to classify submissions. Browsing the database of MR and zb-Math using MSC search is an effective method of the following research in specific areas. Each paper can be involved in multiple study areas with corresponding MSC codes; therefore, it is represented by a combination of various mathematical fields. MSC codes only provide a hierarchical classification scheme; however, it is desirable to apply statistical clustering methods to identify and cluster the specific topics in the academic field by investigating multiple MSC codes presented in the same papers. It is possible to gather detailed information about an intimate relationship between all areas of pure and applied mathematics through a database of mathematical papers.

Due to the emergence of a complex structure of data, various kinds of clustering methods have been studied for partitioning a finite set of data into different clusters according to similarities. Graph clustering algorithms can easily be applied to identify communities when the data is a form of the
Table 1: Top 10 MSC codes most frequently referred

| MSC       | Major category                          | Minor category                        | Num. of docs |
|-----------|-----------------------------------------|----------------------------------------|--------------|
| 01A70     | History and biography                   | Biographies, obituaries, ...           | 23,637       |
| 81T30     | Quantum theory                          | String and superstring theory ...      | 19,333       |
| 65N30     | Numerical analysis                      | Finite elements, Rayleigh-Ritz ...      | 17,567       |
| 68Q25     | Computer science                        | Analysis of algorithms and ...         | 17,187       |
| 47H10     | Operator theory (1959–)                 | Fixed-point theorems ...               | 15,437       |
| 62M10     | Statistics                              | Time series, auto-correlation          | 14,469       |
| 35Q53     | Partial differential equation           | KdV-like equations ...                 | 13,526       |
| 62G05     | Statistics                              | Estimation                             | 13,389       |
| 05C85     | Combinatorics                           | Graph algorithms ...                   | 13,135       |
| 35B40     | Partial differential equation           | Asymptotic behavior of solutions       | 13,014       |

MSC = mathematical subject classification.

network such as social networks or modules in biological systems (Kenley and Cho, 2011). In this paper, we compare graph clustering methods for grouping MSC codes into similar clusters.

This paper is organized as follows. In Section 2, we describe how to obtain the MSC data and construct graph data for applying the graph clustering methods. In Section 3, we review two well-known graph clustering algorithms, entropy-based clustering, and Markov chain clustering (MCL). A comparative study of graph clustering methods is given in Section 4. Concluding remarks are provided in Section 5.

2. Data and preprocessing

In the MR database at the MathSciNet website (http://www.ams.org/mathscinet), each of 3,193,603 entries (papers, books, and reports) provides bibliographic information for a document published from in 1899 to in 2015 along with possibly multiple choices of MSC codes. We download \( P = 2,975,635 \) entries published in 1940–2015 with a least one MSC codes and 1,101,422 documents have a single MSC code, but 62.96% items have multiple codes (995,257 documents with two codes, 531,327 documents with three codes, 228,847 with four codes, and 116,782 documents with five to ten codes).

An MSC code consists of a two digits primary classification code (79 codes are used so far), a letter, and two digits code for minor classification; for example, 00A15 for bibliographies in the 00-General subject. There are 9,395 known MSC codes (as of Jan 1, 2016) but only \( N = 8,822 \) codes are used at least once in the \( P \) items, and we list the top ten MSC codes most frequently referred in Table 1.

We now define the binary incidence matrix \( X = (x_{ip}) \) using discrete MSC code information as:

\[
x_{ip} = \begin{cases} 
1, & \text{if } i^{th} \text{ MSC code is used in an } p^{th} \text{ document}, \\
0, & \text{otherwise,}
\end{cases}
\]

where \( i = 1, \ldots, N \) and \( p = 1, \ldots, P \). Now, we construct a weighted graph \( G = (V, E) \) where the vertices \( V = (v_1, \ldots, v_N) \) and edges \( E \) are the MSC codes and their similarity, respectively. To transform \( X \) into a form of graph, we compute an \( N \times N \) adjacency matrix \( M = XX^T \). Then, the \((i, j)^{th}\) element of \( M, M_{ij} \) represents the strength of edge between vertex \( v_i \) and \( v_j \), i.e., the number of documents containing the \( i^{th} \) and \( j^{th} \) MSC code simultaneously. Intuitively, if \( M_{ij} \) is large, many documents which have both MSC code \( i \) and \( j \), so these MSC codes will be in the same cluster. Table 2 shows 11 edges with edge strength bigger than 2,500; in addition, Figure 1 provides part of the graph that focuses on the field of Statistics (62XXD).

The entries in the adjacency matrix \( M \) associated with the code to document matrix \( X \) represent the
Table 2: Top 11 pairs of MSC codes with $M_{ij}$ larger than 2,500

| MSC$_i$ | Major (Minor) category | MSC$_j$ | Major (Minor) category | $M_{ij}$ |
|---------|------------------------|---------|------------------------|---------|
| 03B35   | Math · · · logic (Mechanization of proofs · · ·) | 68T15   | Computer science (Theorem proving · · ·) | 2,628   |
| 05C85   | Combinatorics (Graph algorithms) | 68R10   | Computer science (Graph theory) | 4,790   |
| 05C85   | Combinatorics (Graph algorithms) | 68Q25   | Computer science (Analysis of algorithm) | 2,801   |
| 17B37   | Nonassociative ring (Quantum groups · · ·) | 81R50   | Quantum theory (Quantum groups · · ·) | 2,972   |
| 35Q30   | Partial differ · · · (Navier-Stokes eq · · ·) | 76D05   | Fluid mechanics (Navier-Stokes eq · · ·) | 3,721   |
| 47H09   | Operator theory (Contraction-type · · ·) | 47H10   | Operator theory (Fixed-point theorem) | 3,097   |
| 47H10   | Operator theory (Fixed-point theorem) | 54H25   | General topology (Fixed-point and · · ·) | 4,688   |
| 60K25   | Probability theory (Queueing theory · · ·) | 90B22   | Operations research (Queues and service) | 5,528   |
| 68Q25   | Computer science (Analysis of algorithm) | 68R10   | Computer science (Graph theory) | 2,949   |
| 81T30   | Quantum theory (String and super · · ·) | 81T60   | Quantum theory (Supersymmetric field) | 2,723   |
| 83E30   | Relativity and gravitation (String and super · · ·) | 83E30   | Relativity and gravitation (String and super · · ·) | 6,412   |

MSC = mathematical subject classification.

Figure 1: A part of graph with the adjacency matrix $M$. 

The number of links between two MSC codes so the contribution to the entries by $p^{th}$ document depends on the number of codes $m_p$ in the document. If a document contains two codes MSC$_i$ and MSC$_j$ then $M_{ij}$ and $M_{ji}$ will be added by one each, so total contribution on the entries of $M$ will be 2. A document with $m_p > 1$ codes make total $P(m_p, 2)$ ($m_p$ permutation 2) contributions to $M$. In order to normalize the contribution to $M$ from each document to be one, we introduce a weighted code-to-document
Figure 2: Relative strength of the edges of $M_{ij}^{w}$.

matrix $X^w$ as:

$$x_{ip}^w = \begin{cases} 
\frac{1}{P(m_p, 2)}, & \text{if } i^{th} \text{ code is one of } m_p \text{ MSC codes in } p^{th} \text{ document}, \\
0, & \text{otherwise,}
\end{cases}$$

for $i = 1, \ldots, N$ and $p = 1, \ldots, P$ with $P(1, 2) := 1$, $P(m, 2) := m(m - 1)$, $m > 1$. Then the sum of all entries of the weighted adjacency $M^w := X^w(X^w)^T$ associated with $X^w$ will be the number of the documents, $\sum_{i=1}^{N} \sum_{j=1}^{N} M_{ij}^{w} = P$.

Table 2 shows the strength of the weighted adjacency matrix. Here for just simple visualization, we only use the graph with 79 nodes associated with the principal classification codes (instead of actual $N = 8,822$ codes for real computations). The width of each edge represents the relative strength of the edges (excluding self-interactions) of $M_{ij}^{w}$. We can easily see that 81-Quantum/83-Relativity, 06-Order.Lattice/03-Math.Logic, 54-Gen.Topology/47-Oper.Theory, and 63-Fluid/35-Part.Diff.Eq pairs are closely related.

3. Methods

In this section, we review two graph clustering methods. Consider the clustering problem for partitioning $N$ vectors of $P$-dimensional binary data $(X)$ or vertices $V$ of $G$ into $K(\leq N)$ disjoint clusters $C_1, C_2, \ldots, C_K$. Let us define a clustering set $C$ as a collection of $K$ disjoint clusters $C_1, C_2, \ldots, C_K$. Note that a clustering set $C$ of $K$ disjoint clusters also can be regarded as an $K$-partition of $N=[1, 2, \ldots, N]$. Then, $i \in C_k$ implies that an $i^{th}$ MSC code belongs to $k^{th}$ cluster $C_k$. 
3.1. Entropy-based clustering algorithm

Entropy type measures for similarity among the finite set of data have been frequently used. Since it is hard to define a distance when the given set of data is discrete, we use the entropy-type measures for similarity among distinct objects. Here, we review an entropy-based clustering method Li et al. (2004), Chen and Liu (2005) minimizing the expected entropy of the partition. An entropy-based clustering algorithm can be formally derived in the framework of probabilistic clustering models.

For given cluster \( C_k \), an entropy of the partition is defined as

\[
H(C_k) = - \sum_{i \in C_k} \sum_{s=0}^{P} \Pr(x_{ip} = s) \log \Pr(x_{ip} = s),
\]

for \( k = 1, \ldots, K \). For each cluster \( C_k \), we let \( n_k \) be the number of MSC codes in each cluster \( C_k \) and \( N_{p,k,1} = \sum_{i \in C_k} x_{ip}, N_{p,k,0} = n_k - N_{p,k,1} \). Then, with the estimated \( \Pr(x_{ip} = s) = \hat{\Pr}(x_{ip} = s) = N_{p,k,s}/n_k \), the estimate of \( H(C_k) \) is

\[
\hat{H}(C_k) = - \sum_{p=1}^{P} \sum_{s=0}^{1} \frac{N_{p,k,s}}{n_k} \log \frac{N_{p,k,s}}{n_k}.
\]

An entropy-based clustering method finds \( C_1, \ldots, C_K \) to minimize the following entropy measure of the partition

\[
\hat{H}(C) = \frac{1}{N} \sum_{k=1}^{K} n_k \hat{H}(C_k)
\]

which is the weighted sum of an estimated entropy \( \hat{H}(C_k) \).

Following Li et al. (2004), \( \hat{H}(C) \) can be approximated as

\[
\hat{H}(C) \approx \frac{1}{N} \sum_{k=1}^{K} \frac{1}{n_k} \sum_{i,j \in C_k} \sum_{p=1}^{P} |x_{ip} - x_{jp}|
\]

\[
= \frac{1}{N} \sum_{k=1}^{K} \frac{1}{n_k} \sum_{i,j \in C_k} \left( x_{iv} + x_{jv} - M_{ij} \right),
\]

where \( x_{iv} = \sum_{p=1}^{P} x_{ip} \) and \( x_{jv} = \sum_{p=1}^{P} x_{jp} \). Note that to minimize the estimated entropy, the pairs of MSC codes with large \( M_{ij} \) tend to be in the same cluster.

For an actual algorithm implementation, we use the Monte-Carlo method suggested by Li et al. (2004). First, we make all MSC codes placed randomly in the \( K \) clusters with almost the same size. Then we choose one MSC code \( v \) randomly. Let \( C_i \) denote a current cluster containing a randomly chosen code \( v \). According to the value of an estimated expected entropy, we decide if the randomly chosen code \( v \) moves from a current cluster \( C_i \) to another cluster \( C_j \) (\( \neq C_i \)) or not. We move the code \( v \) from a current cluster \( C_i \) to another cluster \( C_j \) if the cluster \( C_j \) (after getting a current code \( v \)) has less estimated expected entropy than a current cluster \( C_i \). Otherwise, the code \( v \) stays in the current cluster \( C_i \). After continuing these steps, we assign MSC codes into \( K \) disjoint clusters minimizing the estimated expected entropy. The repetitive relocation is required in the actual algorithm implementation. Note that the convergence property of the Monte-Carlo optimization is shown in Li et al. (2004).
3.2. Markov chain clustering algorithm

Recall that the \((i, j)\)th element of an \(N \times N\) adjacency matrix \(M\), \(M_{ij}\) represents the number of papers containing the \(i\)th and \(j\)th MSC code. Define a transition matrix \(T\) for the graph \(G\) as a matrix having an \((i, j)\)th element \(T_{ij}\) which is obtained as

\[
T_{ij} = \frac{M_{ij}}{\sum_{j \neq i} M_{ij}},
\]

where \(1 \leq i \leq N\) and \(1 \leq j \leq N\).

A random walk of a directed graph consists of a series of vertices generated at a starting vertex, reaching the next vertex by selecting one of the outgoing edges, and repeating that process. A finite Markov chain defined on the set of states, which is denoted by \(V = \{v_1, v_2, \ldots, v_N\}\), is specified by the transition matrix \(T\) of the weighted graph \(G\) on the vertex set \(V\). Therefore, there is no significant difference between a random walk on graphs and a finite Markov chain. In the weighted directed graph, all Markov chains can be thought in random walks on directed edges. A MCL algorithm generates random walks on weighted graphs for dividing a finite set of data into different clusters according to their similarity. When a random walk reaches the vertex \(w\) from a starting vertex \(v\) with high probability, vertices \(v\) and \(w\) need to be collected together into the same cluster in the MCL algorithm.

Let that an \((i, j)\)th element of the matrix \(M^r\), which is made by the \(r\)th power of \(M\), is a number of length \(r\) paths from \(v_i\) to \(v_j\). Therefore a number of length \(r\) paths between two vertices, \(v_i\) and \(v_j\), is obtained by raising the adjacency matrix of \(G\) to the exponent \(r\). To subdivide MSC codes into \(K\) disjoint clusters according to their similarity, the MCL algorithm is explained in the following four steps (Van Dongen, 2008).

- **Step 1.** We define a matrix \(M^*\) having an \((i, j)\)th element which is obtained as

\[
M^*_{ij} = \begin{cases} 
M_{ij}, & \text{when } i \neq j, \\
1, & \text{when } i = j.
\end{cases}
\]

Then we execute a normalization of the matrix \(M^*\). Let us call the normalized matrix \(M\).

- **Step 2 (Expansion).** In this step, we execute the power of the matrix \(M\), which means \(M \times M\). Then we again call the powered matrix \(M\). Note that the effectiveness of the power of a matrix diminishes as the flow proceeds, where a flow on the graph \(G\) is defined as a transition probability from a vertex to another vertex. The higher powers of the matrix obtained after repeating this step also reinforce more connections between MSC codes.

- **Step 3 (Inflation).** Let \(r\) be a given inflation parameter. Depending on the given inflation parameter \(r\), an inflation operator \(\Gamma_r : \mathbb{R}^{N \times N} \rightarrow \mathbb{R}^{N \times N}\) is defined as

\[
\Gamma_r M_{ij} = \frac{(M_{ij})^r}{\sum_{q=1}^{N} (M_{qj})^r},
\]

where \(\Gamma_r M_{ij}\) is an \((i, j)\)th element of the matrix \(\Gamma_r M\).

Note that the inequality of each column becomes bigger after taking an inflation operator \(\Gamma_r\). If an \((i, j)\)th element of the matrix \(M\) has a big (small) value, then an \((i, j)\)th element of the matrix
\( \Gamma, M \) becomes bigger (smaller), respectively. After executing a normalization of the matrix \( \Gamma, M \), we again call the normalized matrix \( M \).

• Step 4 (Iteration). We repeat the second and third steps until the matrix \( M \) converges.

The expansion step strengthens more connections among MSC codes and subdivides MSC codes into some vast clusters. However, the inflation step both strengthens already strong currents and weakens already weak currents. To avoid dominantly massive clusters, we need the inflation step that distributes the MSC codes across the different clusters. In the MCL algorithm, the given inflation parameter \( r \) controls the extent of strengthening and weakening and influences the granularity of disjoint clusters. For the given set of data, we subdivide MSC codes into 136 disjoint clusters by setting a value 1.5 as an inflation parameter.

4. Numerical study

In this section, we apply the entropy-based clustering (Entropy) in Section 3.1 and MCL in Section 3.2 to the MSC code dataset. Furthermore, we compare these with conventional methods and evaluate the performance. As the traditional methods, we use the hierarchical clustering algorithm. First, we start by assigning each MSC code into a single cluster. Then we have \( N \) clusters containing just one MSC code. Second, we find the closest pair of clusters, which is most similar among them. Then we merge them into a single cluster. Now we have \( N - 1 \) clusters. Third, we compute all distances between any two clusters among \( N - 1 \) clusters, where a distance between two clusters, \( C_i \) and \( C_j \), is defined as an average of all distances between any two codes \( v \in C_i \) and \( w \in C_j \). Then we repeat the second step to find the closest pair of clusters that are most similar and merge them into a single cluster. We repeat the second and third steps until we get \( K \) disjoint clusters.

For dissimilarity measure \( \delta(x_i, x_j) \) between two MSC codes \( i \) and \( j \), we use the following Jaccard coefficient and Simple matching coefficient:

\[
\text{(1) Jaccard coefficient:} \quad \sum_{p=1}^{P} \frac{|x_{ip} - x_{jp}|}{\|\{p : x_{ip} = 1 \text{ or } x_{jp} = 1\}\|}.
\]

\[
\text{(2) Simple matching coefficient:} \quad \sum_{p=1}^{P} \frac{|x_{ip} - x_{jp}|}{P}.
\]

In this paper, the hierarchical clustering algorithms based on Jaccard coefficient and Simple matching coefficient are called hclust1 and hclust2. The hclust1 method considers the effectiveness of the number of the coordinates of two codes, \( x_i = (x_{i1}, x_{i2}, \ldots, x_{ip}) \) and \( x_j = (x_{j1}, x_{j2}, \ldots, x_{jp}) \), having the value 1 in the same position.

We consider within-cluster criterion for the clustering set \( C \) that evaluates the clustering results of these four algorithms (Entropy, hclust1, hclust2, and MCL)

\[
D(C) = \frac{1}{N} \sum_{k=1}^{K} \frac{1}{n_k} \sum_{i,j \in C_k} \delta(x_i, x_j),
\]

where a dissimilarity \( \delta(x_i, x_j) \) is a Jaccard coefficient or a Simple matching coefficient between two MSC codes, \( x_i \) and \( x_j \). The within-cluster criteria for the clustering set \( C \) of Entropy, hclust1, hclust2, and MCL are given in Table 3. All clustering methods are similar in terms of the Jaccard coefficient.
Table 3: Dissimilarity measures of clustering methods

|               | MCL     | Entropy | hclust1   | hclust2   |
|---------------|---------|---------|-----------|-----------|
| Jaccard       | 0.982   | 0.983   | 0.984     | 0.984     |
| Simple matching coefficient | 4.59 \times 10^{-4} | 4.39 \times 10^{-4} | 4.68 \times 10^{-4} | 3.81 \times 10^{-4} |

MCL = Markov chain clustering; Entropy = entropy-based clustering; hclust1 = hierarchical clustering algorithms based on Jaccard coefficient; hclust2 = hierarchical clustering algorithms based on Simple matching coefficient.

Table 4: Size of biggest four clusters

|          | Entropy | hclust1 | hclust2 | MCL     |
|----------|---------|---------|---------|---------|
| Biggest  | 926     | 6278    | 8666    | 601     |
| Second   | 634     | 1605    | 4469    |         |
| Third    | 526     | 792     | 3       | 423     |
| Fourth   | 398     | 8       | 2       | 395     |
| Standard deviation | 115.3391 | 556.0134 | 740.2677 | 96.46783 |

Entropy = entropy-based clustering; hclust1 = hierarchical clustering algorithms based on Jaccard coefficient; hclust2 = hierarchical clustering algorithms based on Simple matching coefficient; MCL = Markov chain clustering.

Table 5: Estimated entropy and W-Ratio value for each clustering method

| Algorithm | Estimated entropy | W-Ratio |
|-----------|-------------------|---------|
| MCL       | 3832.066          | 0.71434 |
| Entropy   | 2604.859          | 0.45628 |
| hclust1   | 6094.500          | 0.94477 |
| hclust2   | 5196.884          | 0.77836 |

MCL = Markov chain clustering; Entropy = entropy-based clustering; hclust1 = hierarchical clustering algorithms based on Jaccard coefficient; hclust2 = hierarchical clustering algorithms based on Simple matching coefficient.

In this data, 97.8 percent of MSC code pairs has a value 1 as a Jaccard coefficient. This means that the number of MSC code pairs used in different papers is extremely small. hclust2 shows the smallest simple matching coefficient; however, hclust1 and hclust2 insert most MSC codes into one or two big clusters (Table 4).

Next, we calculate the estimated entropy value $\hat{H}(C)$ (3.1) with the disjoint clusters $C_1, \ldots, C_K$ of each clustering algorithm and the W-Ratio as another evaluation measure. Let $T = \{T_1, T_2, \ldots, T_P\}$ be the set of documents (papers) where an $p^{th}$ paper $T_p$ contains $m_p$ MSC codes, where $1 \leq p \leq P$. For given clustering set $C$, W-Ratio is defined as

$$W(C) = \frac{\sum_{k=1}^{K} \sum_{i,j \in C_k} \sum_{p=1}^{P} W_p(i, j)}{\sum_{i,j \in C} \sum_{p=1}^{P} W_p(i, j)},$$

where

$$W_p(i, j) = \begin{cases} 
1/m_p(m_p-1), & \text{if an } p^{th} \text{ document } T_p \text{ contains distinct codes } i \text{ and } j, \\
0, & \text{if } i = j \text{ or an } p^{th} \text{ document } T_p \text{ does not contain codes } i \text{ or } j.
\end{cases}$$

Note that to get relatively high W-Ratio, the pairs of MSC code with the large number of papers containing them tend to be in the same cluster.

Table 5 shows the estimated entropy value and W-Ratio value of four algorithms. The estimated entropy of the Markov Chain clustering (MCL) algorithm is 3832.066, and hclust1 and hclust2 have larger entropies. The entropy-based algorithm, the expected entropy value of $K$ disjoint clusters is
Graph clustering of MSC

Figure 3: Estimated entropy values of graph clustering algorithms. MCL = Markov chain clustering; Entropy = entropy-based clustering; hclust1 = hierarchical clustering algorithms based on Jaccard coefficient; hclust2 = hierarchical clustering algorithms based on Simple matching coefficient.

decreasing on iterations and we can see that it converges to the smallest amount, 2604.859 (Figure 3). MCL and hclust1 algorithms have relatively high $W$-Ratio values. Note the entropy-based algorithm that makes clustering groups relatively uniform, which does not consider the number of papers containing the same MSC codes, has a relatively low $W$-Ratio value. However, the hclust2 algorithm having one or two dominantly big clusters has a relatively high $W$-Ratio value despite not considering the effectiveness of the number of papers containing the same MSC codes.

5. Concluding remarks

In this paper, we compare appropriate two graph clustering methods for grouping Mathematical Subject Classification (MSC) codes according to similarities. The hclust1 and hclust2 clustering algorithms have one or two dominantly big clusters; therefore, we compare the Markov chain and entropy-based clustering algorithms making clustering groups relatively uniform. The MCL clustering algorithm has a higher entropy expectation value than the entropy-based clustering algorithm minimizing the expected entropy of the partition. The MCL clustering algorithm has a larger $W$-Ratio value than the entropy-based clustering algorithm since the MCL algorithm considers the number of papers containing the same MSC codes; however, the entropy-based clustering algorithm does not. Therefore, we conclude that the MCL clustering algorithm is suitable for clustering the MSC codes.

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