Lipschitz Functions on Expanders are Typically Flat

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This work studies the typical behaviour of random integer-valued Lipschitz functions on expander graphs with sufficiently good expansion. We consider two families of functions: $M$-Lipschitz functions (functions which change by at most $M$ along edges) and integer-homomorphisms (functions which change by exactly 1 along edges). We prove that such functions typically exhibit very small fluctuations. For instance, we show that a uniformly chosen $M$-Lipschitz function takes only $M+1$ values on most of the graph, with a double exponential decay for the probability of taking other values.

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1. Introduction

In this work we investigate the typical behaviour of random Lipschitz functions on expander graphs. We focus on the following two models. An $M$-Lipschitz function on a graph $G$ is an integer-valued function on the vertices of $G$ which changes by at most $M$ between adjacent vertices. Similarly, a $\mathbb{Z}$-homomorphism (or simply a homomorphism) on $G$ is an integer-valued function on the vertices of $G$ which changes by exactly one between adjacent vertices. Taking the graph $G$ to be a finite expander graph, we consider

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the typical properties of functions chosen uniformly at random from one of these families (fixing the function value to be zero at some fixed vertex).

One motivation for this work comes from previous investigations of \( \mathbb{Z} \)-homomorphisms on several tree-like graphs [1], on the hypercube [7, 10] and on (finite boxes in) the lattice \( \mathbb{Z}^d \) for large \( d \) [11]. These suggest that typical Lipschitz functions on highly connected graphs tend to exhibit very small fluctuations. Indeed, such behaviour can also be expected from a comparison with random surface models in statistical mechanics, such as the Gaussian free field (see also [2]). Expander graphs are natural candidates to test this paradigm, as they are highly connected graphs which are important for many applications; see the survey [8]. It is also well known that most graphs are expanders.

An additional motivation is to try and understand to what extent the local behaviour of typical Lipschitz functions on a graph is affected by the global features of the graph. Many expander graphs locally have the structure of a tree, though globally they are very far from being a tree. Will the typical Lipschitz function on an expander graph exhibit locally large fluctuations, as it does on a tree, or will it exhibit locally small fluctuations, as suggested by the global structure?

Our results apply to expander graphs with sufficiently good expansion, in a certain quantitative sense (see Definitions III and VI). We show that on such graphs, the typical Lipschitz function exhibits very small fluctuations. Thus, although the graph may look like a tree locally, it is the global structure of it which determines the local fluctuations. More specifically, we find that a random \( M \)-Lipschitz function will take only \( M + 1 \) different values on most of the graph. Moreover, the probability that at any fixed vertex the function takes a value which is further than \( tM \) from this set of \( M + 1 \) values decays like a double exponential in \( t \). As a result, the maximum value of the function is, with high probability, of order \( \log \log n \), where \( n \) is the number of vertices of the graph.\(^1\)

Similar results are obtained for \( \mathbb{Z} \)-homomorphisms, where it is required that the underlying graph is bipartite. There, we find that the typical function takes predominantly one value on one of the colour classes and two values on the other colour class, again with a double exponential decay of the probability of taking other values. In a somewhat different setting, our methods also yield that grounded Lipschitz functions on \( d \)-ary trees, \textit{i.e.}, functions constrained to take the value zero on all leaves of the tree, exhibit similarly small fluctuations.

While our methods require the graph to have sufficiently good expansion, it is not clear to what extent this requirement is necessary. A discussion with several open questions is presented in Section 6.

1.1. \( M \)-Lipschitz functions

\( M \)-Lipschitz functions on graphs are defined as follows.

\begin{definition}[\( M \)-Lipschitz functions] Let \( v_0 \) be a fixed vertex in a graph \( G \) and let \( M \) be a positive integer. Denote by \( \text{Lip}_{v_0}(G; M) \) the family of \( M \)-Lipschitz functions from the

\(^1\) Logarithms in this text are of base \( e \).
vertex set of $G$ to $\mathbb{Z}$ that send $v_0$ to the origin, that is, the family of maps $f : V(G) \to \mathbb{Z}$ such that $f(v_0) = 0$ and $|f(u) - f(v)| \leq M$ for every $\{u, v\} \in E(G)$.

We note for later use that if $G$ is a finite connected graph, then $\text{Lip}_{v_0}(G; M)$ is a finite set.

There are several equivalent definitions of expander graphs. Our definition is inspired by the so-called expander mixing lemma (see [8]), which relates the edge distribution of a regular graph to the spectral properties of its adjacency matrix. For two subsets $A, B$ of the vertices of a graph $G$, denote by $E(A, B)$ the set of pairs $(a, b) \in A \times B$ such that $\{a, b\}$ is an edge of $G$ and denote $e(A, B) = |E(A, B)|$.

**Definition II (expander).** A $d$-regular $n$-vertex graph $G$ is called a $\lambda$-expander if, for all $S, T \subseteq V(G)$, we have

$$\left| e(S, T) - \frac{d}{n}|S||T| \right| \leq \lambda \sqrt{|S||T|}.$$  

Every $d$-regular graph is an expander with $\lambda = d$ and hence the definition becomes meaningful only for $\lambda < d$. Simple examples show that expanders cannot be too good: for $S = \{v\}$ and $T$ the $d$ neighbours of $v$ in $G$ we obtain $\lambda \geq \sqrt{d(1 - d/n)}$ and for $S = T = \{v\}$, we have $\lambda \geq d/n$ so that in particular, $\lambda \geq 1/2$.

In our theorems, we consider sufficiently good expanders in the sense that $\lambda$ is required to be smaller than some specific function of $d$ and $M$. It can be shown that if $d$ is sufficiently large, then most $d$-regular graphs are good expanders. For example, Friedman showed that for every $d > 2$ and every $\varepsilon > 0$, with probability tending to 1 as $n$ tends to infinity, a uniformly chosen random $n$-vertex $d$-regular graph is a $(2\sqrt{d - 1} + \varepsilon)$-expander [6].

**Definition III (M-good expander).** A graph $G$ is an $M$-good expander if $G$ is a $d$-regular $\lambda$-expander with

$$\lambda \leq \frac{d}{32(M + 1) \log(9Md^2)}.$$  

Our main result shows that a typical $M$-Lipschitz function on an $M$-good expander is locally very flat. In particular, such a function takes values in a set of $M + 1$ consecutive integers at all but an exponentially small fraction of the vertices, where exponentially small is with respect to the parameters of the expander. The first step towards establishing this property is to note that every $M$-Lipschitz function takes values in a set of $M + 1$ consecutive integers at all but a polynomially small fraction of the vertices.

**Lemma 1.1.** Let $G$ be an $n$-vertex $d$-regular $\lambda$-expander and $v_0 \in V(G)$. For every $f \in \text{Lip}_{v_0}(G; M)$, there exists $k \in \mathbb{Z}$ such that

$$|\{v : f(v) \notin \{k, k + 1, \ldots, k + M\}\}| \leq \frac{2\lambda n}{d}.$$  

(1.1)
Moreover, there is a way to associate to each \( f \in \text{Lip}_{v_0}(G; M) \) which is not identically zero an interval of the form \( \{k, k+1, \ldots, k+M\} \) satisfying (1.1), which we denote by \( \text{phase}(f) \), in such a manner that \( \text{phase}(-f) = -\text{phase}(f) \).

The lemma allows us to define \( \text{phase}(f) \) for \( M \)-Lipschitz functions which are not identically zero. For completeness, we define the phase of the zero function as \( \{0\} \). The lemma does not provide information about the range of values which the function can take. Indeed, the diameter of an \( n \)-vertex expander graph \( G \) is of the order of \( \log n \) (see also Corollary 2.6 below) and hence there exist Lipschitz functions on \( G \) taking order \( \log n \) distinct values (though they still must satisfy Lemma 1.1). The next theorem and corollary, which are our main results, show that such large fluctuations are highly atypical.

To discuss the typical properties of an \( M \)-Lipschitz function, we consider the uniform probability measure on \( \text{Lip}_{v_0}(G; M) \). Restricting to finite connected graphs \( G \), we denote by \( f \in \mathbb{R} \text{Lip}_{v_0}(G; M) \) a uniform random element in \( \text{Lip}_{v_0}(G; M) \). To describe our results, we use the metric of \( G \) (the distance between a pair \((u, v)\) of vertices is the length of a shortest path connecting \( u \) and \( v \) in \( G \)), and we denote by \( B(v, t) \) the ball of radius \( t \) around a vertex \( v \). Finally, for an integer \( b \) and a non-empty set \( A \subseteq \mathbb{Z} \), define \( \text{dist}(b, A) = \min_{a \in A} |b - a| \).

**Theorem 1.2.** Let \( G \) be a connected \( M \)-good expander and \( v_0 \in V(G) \). Let \( f \in \mathbb{R} \text{Lip}_{v_0}(G; M) \). For every positive integer \( t \) and every \( v \in V(G) \),

\[
\mathbb{P}(\text{dist}(f(v), \text{phase}(f)) > (t - 1)M) \leq \exp\left(-\frac{|B(v, t)|}{5(M + 1)}\right).
\]  

(1.2)

Let us make a few comments.

(1) Proposition 2.5 below shows that in any expander, \( |B(v, t)| \) grows at least exponentially in \( t \). Thus, under the conditions of Theorem 1.2, we have

\[
\mathbb{P}(\text{dist}(f(v), \text{phase}(f)) > (t - 1)M) \leq \exp(-Ct)
\]

for some \( C = C(d, \lambda) > 1 \); the value of \( C \) can be calculated from (1.2) and (2.2). If the girth of \( G \) is at least \( 2t + 1 \), we have \( |B(v, t)| \geq d(d - 1)^{t-1} \), which can improve the estimate further.

(2) If we apply Theorem 1.2 with \( v = v_0 \), we obtain bounds on the probability distribution of \( \text{phase}(f) \).

(3) As a simple consequence of Theorem 1.2 and Proposition 2.5, we see that the variance of \( f(v) \) is uniformly bounded for all \( v \in V(G) \) by \( C_1 M^2 \), where \( C_1 \) is an absolute constant.

We can use Theorem 1.2 and Proposition 2.5 to obtain a bound on the maximum of \( f \) (or alternatively, on the range of values \( f \) takes). By applying a union bound over all vertices of \( G \), we obtain the following estimate.
Corollary 1.3. Let $G$ be a connected $M$-good expander and $v_0 \in V(G)$. Let $f \in \text{Lip}_{v_0}(G;M)$. There exists an absolute constant $C$ such that

$$\mathbb{P}\left(\max(f) > \frac{C}{\log(d/\lambda)} M \log(n)\right) \leq \frac{1}{n^2}.$$ 

In particular, $\mathbb{E}[\max(f)] \leq \frac{2C}{\log(d/\lambda)} M \log(n)$.

It appears that by applying the techniques of Benjamini, Yadin and Yehudayoff [3] one may obtain a converse to the above corollary, showing that $\mathbb{E}[\max(f)] \geq c M \log(n)$ for some $c$ depending only on $d$ and $\lambda$. We do not explore this direction here, but mention that the general approach is to consider the possibility that $f$ grows quickly in a ball of small radius and then use the fact that there are many such disjoint balls.

1.2. Homomorphisms

We also study the typical behaviour of graph homomorphisms from $G$ to $\mathbb{Z}$ or, as we will refer to them below, homomorphism height functions or simply homomorphisms.

Definition IV (homomorphisms). Let $v_0$ be a fixed vertex in the graph $G$. Let $\text{Hom}_{v_0}(G)$ denote the family of graph homomorphisms from $G$ to $\mathbb{Z}$ that send $v_0$ to the origin, that is, the family of maps $f : V(G) \to \mathbb{Z}$ so that $f(v_0) = 0$ and $|f(u) - f(v)| = 1$ for every $\{u,v\} \in E(G)$.

The family $\text{Hom}_{v_0}(G)$ is non-empty if and only if $G$ is bipartite. It is finite if and only if $G$ is finite, connected and bipartite. Let $G$ be bipartite with colour classes $V_0$ and $V_1$ and assume that $v_0$ is in $V_0$. Then, every $f \in \text{Hom}_{v_0}(G)$ takes even values on vertices of $V_0$ and odd values on $V_1$.

Since bipartite graphs contain very large independent sets, they are clearly not expanders in the sense of Definition II. Therefore, we shall consider the following (standard) bipartite analogue of expander graphs.

Definition V (bi-expander). A $d$-regular bipartite graph $G$ with $2n$ vertices\(^2\) is called a $\lambda$-bi-expander if the following holds. Let $V_0$ and $V_1$ be the two colour classes of $G$. For all $S \subseteq V_0$ and $T \subseteq V_1$,

$$|e(S, T) - \frac{d}{n} |S||T| | \leq \lambda \sqrt{|S||T|}.$$ 

As in the case of (non-bipartite) $\lambda$-expanders, every $d$-regular bipartite graph is a $\lambda$-bi-expander with $\lambda = d$ and hence Definition V becomes meaningful only for $\lambda < d$. In our theorems, we consider sufficiently good expanders, in the sense that $\lambda$ is required to be

\(^2\)Since every regular bipartite graph has an even number of vertices, we choose to denote this number by $2n$. This is somewhat inconsistent with our notation from Section 1.1, but will make later formulas look considerably cleaner.
smaller than some function of $d$. As in the non-bipartite case, it can be shown that most $d$-regular bipartite graphs are good bi-expanders, provided that $d$ is sufficiently large.

**Definition VI (good bi-expander).** A graph $G$ is a good bi-expander if $G$ is a $d$-regular $\lambda$-bi-expander with $\lambda \leq \frac{d}{300 \log d}$.

We will show that homomorphism height functions on good bi-expanders are very flat in the sense that a typical function is nearly constant on one of the two colour classes. Here, nearly constant will mean constant on all but an exponentially (in the parameters of the expander) small fraction of the vertices. As in the $M$-Lipschitz case, we first note that every homomorphism height function is constant on one of the two colour classes, apart from a polynomially small fraction of the vertices.

**Lemma 1.4.** Let $G$ be a $2n$-vertex $d$-regular $\lambda$-bi-expander and $v_0 \in V(G)$. For every $f \in \text{Hom}_{v_0}(G)$, there exist $i \in \{0, 1\}$ and $k \in \mathbb{Z}$ such that

$$|\{v \in V_i : f(v) \neq k\}| \leq \frac{2\lambda n}{d}. \quad (1.3)$$

Let $i^* \in \{0, 1\}$ be the smallest $i$ so that there exists a $k$ satisfying (1.3) with $V_i$. Denote by $\text{phase}(f)$ the smallest $k$ satisfying (1.3) for $i^*$. The parity of $\text{phase}(f)$ is $i^*$. Moreover, if $\lambda < d/3$, then

$$|\{v \in V(G) : |f(v) - \text{phase}(f)| \geq 2\}| \leq \frac{3\lambda n}{d}. \quad (1.4)$$

As before, we restrict our attention to finite, connected, bipartite $G$, and let $f \in_R \text{Hom}_{v_0}(G)$ denote a uniform random element in $\text{Hom}_{v_0}(G)$. The following theorem is our main result for homomorphism height functions.

**Theorem 1.5.** Let $G$ be a connected good bi-expander and $v_0 \in V(G)$. Let $f \in_R \text{Hom}_{v_0}(G)$. For every integer $t \geq 1$ and every $v \in V(G)$,

$$\mathbb{P}(|f(v) - \text{phase}(f)| > t) \leq \exp\left(-\frac{|B(v, t)|}{3}\right). \quad (1.5)$$

Comments similar to those after Theorem 1.2 apply here as well. Proposition 2.10 shows that $|B(v, t)|$ grows at least exponentially in $t$ and hence, under the assumptions of the theorem,

$$\mathbb{P}(|f(v) - \text{phase}(f)| > t) \leq \exp(-C't)$$

for some $C = C(d, \lambda) > 1$; the value of $C$ can be calculated from (1.5) and (2.4). Further, by taking $v = v_0$ we may obtain bounds on the distribution of $\text{phase}(f)$, and so the theorem implies that $\max_{v \in V(G)} \text{Var} f(v)$ is bounded above by an absolute constant.

Observing that the diameter of $2n$-vertex bi-expanders is logarithmic in $n$ (see Corollary 2.11), we see that every homomorphism height function can take at most order
log \( n \) values. Our next statement sharpens this bound considerably by showing that the typical order of magnitude of the range of the random homomorphism height function is \( \log(\log n) \). The lower bound in this theorem follows from the results of [3].

**Theorem 1.6.** Let \( G \) be a good 2n-vertex bi-expander and let \( v_0 \in V(G) \). There exist a constant \( c > 0 \), which depends only on \( d \) and \( \lambda \), and an absolute constant \( C \) such that if \( f \in \mathbb{R} \text{Hom}_{v_0}(G) \), then

\[
\Pr\left(c \log(\log n) \leq \max(|f|) \leq \frac{C}{\log(d/\lambda)} \log(\log n)\right) \geq 1 - \frac{1}{n^2}.
\]

1.3. Grounded Lipschitz functions on trees

As a final example of the applicability of our methods, we study Lipschitz functions on trees. We denote by \( T^d_h \) the complete \((d - 1)\)-ary tree of height \( h \geq 1 \) (in which all internal vertices, including the root, have degree \( d \)), denote its root vertex by \( v_r \) and the set of its leaves by \( V_L \). By definition, \( T^d_h \) is a tree with \( d(d - 1)^{h-1} \) leaves, all at distance exactly \( h \) from \( v_r \). It is a simple observation that a uniformly chosen function from \( \text{Hom}_{v_r}(T^d_h) \) will behave like a random walk along branches of the tree, and thus typically take values of order \( \sqrt{h} \) at the leaves of the tree and reach a value of order \( Mh \) at its maximum. Similarly, a uniformly chosen function from \( \text{Lip}_{v_r}(T^d_h; M) \) typically takes values of order \( M\sqrt{h} \) at the leaves of the tree and reaches a value of order \( Mh \) at its maximum.

We consider instead a different probabilistic model. We let \( \text{Lip}_{V_L}(T^d_h; M) \) be the set of \( M\)-Lipschitz functions on \( T^d_h \) which equal zero on \( V_L \), and similarly let \( \text{Hom}_{V_L}(T^d_h) \) be the set of homomorphism height functions on \( T^d_h \) which equal zero on \( V_L \). We call such functions *grounded*. Benjamini, Häggström and Mossel [1] investigated typical grounded homomorphism functions on \( T^d_h \) and showed that they exhibit very different behaviour from homomorphism functions normalized at the root of the tree. For example, they showed that for any \( d > 2 \), if \( f \in \mathbb{R} \text{Hom}_{V_L}(T^d_h) \), then

\[
\Pr(|f(v_r)| > t) \leq 2e^{ct},
\]

where \( c = c(d) \leq 2^{-d+1} \), so that, in particular, the distribution of \( f(v_r) \) is tight as \( h \to \infty \).

Using the same techniques as for expander graphs, we establish similar results for grounded \( M\)-Lipschitz functions on trees. Also, in the next section, we explain our argument briefly in the context of grounded homomorphism functions. Our method seems simpler than that of [1] but has the disadvantage that it does not apply for small \( d \).

**Theorem 1.7.** For any \( M \geq 1 \), \( d > 40(M + 1)\log(M + 1) \) and \( h \geq 1 \), if \( f \in \mathbb{R} \text{Lip}_{V_L}(T^d_h; M) \), then

\[
\Pr(|f(v)| > tM) \leq \exp\left(-\frac{d(d - 1)^{h-1}}{5(M + 1)}\right),
\]

for all vertices \( v \) of \( T^d_h \) and integers \( t \geq 1 \).

In Section 6 we discuss extensions of the above theorem to any \( d > 2 \) and \( M \geq 1 \).
1.4. Proof idea

We now briefly discuss the general idea behind the proofs. Our goal is to show that a certain set is small (the set of functions \( f \) so that \( f(v) \) is far from \( \text{phase}(f) \)). The basic tool we use is extremely simple. By definition, a set \( P \) is not larger than another set \( Q \) if there exists a one-to-one map from \( P \) to \( Q \). To show that \( P \) is much smaller than \( Q \), we just need to find a few-to-many map \( T \) from \( P \) to \( Q \), as the following (trivial) lemma shows.

**Lemma 1.8.** Let \( P \) and \( Q \) be finite sets. Let \( T \) be a map on \( P \) such that \( T(p) \subseteq Q \) for all \( p \in P \). If there are \( \alpha, \beta > 0 \) such that for each \( p \in P \) and each \( q \in Q \),

\[
|T(p)| \geq \beta \quad \text{and} \quad |\{p' \in P : q \in T(p')\}| \leq \alpha,
\]

then \(|P|/|Q| \leq \alpha/\beta|.

**Proof.** We have \( \beta|P| \leq \sum_{p \in P} |T(p)| = \sum_{q \in Q} |\{p' \in P : q \in T(p')\}| \leq \alpha|Q| \).

The heart of the proof is constructing such a map \( T \) and establishing that it is few-to-many. Let us focus on the specific case of grounded homomorphisms on trees since it is the simplest case to explain and already captures the essence of our approach. Let \( G = T^d_h \) be the complete \((d-1)\)-ary tree of height \( h \geq 1 \) (in which all internal vertices have degree \( d \)). Recall that a grounded homomorphism on \( G \) is a function \( f : V(G) \to \mathbb{Z} \) changing by exactly one along every edge and having \( f(v_0) = 0 \) for every leaf \( v_0 \). The goal is to estimate the probability that a random grounded homomorphism takes a large value at some fixed vertex \( v \) of \( G \). More precisely, for some integer \( t \geq 1 \), we want to estimate \( P(f(v) > t) \). By symmetry, this is the same as \( P(f(v) < -t) \). An interesting case to keep in mind is the case when \( v \) is the root vertex. While it is possible for a grounded homomorphism function to take a large value at the root, we will show that this behaviour is highly atypical.

We denote by \( Q \) the set of all grounded homomorphisms on \( G \) and by \( P \) the set \( \{f \in Q : f(v) > t\} \). We wish to show that the size of \( P \) is much smaller than the size of \( Q \) and we will do so by constructing a few-to-many map \( T \) from \( P \) to \( Q \) and invoking Lemma 1.8. For fixed \( f \) in \( P \), we now describe how to define \( T(f) \).

Let \( G^{\leq 2} \) be the graph defined on the same vertex set as \( G \) so that two vertices are connected in \( G^{\leq 2} \) if their distance in \( G \) is at most 2. Consider \( A = A(f) \), where

\[
A(f) \text{ is the connected component of } v \text{ in the subgraph of } G^{\leq 2} \text{ induced by the set of vertices } w \text{ having } f(w) \geq 2.
\]

Let \( X = X(f) \) be the outer boundary of \( A \) in \( G \), i.e., the set of vertices \( w \) which are not in \( A \) but are connected to \( A \) by an edge. Let \( Y = Y(f) \) be the outer boundary of \( A \cup X \) in \( G \). So, we have a set \( A \) and two ‘shells’ of it, \( X \) and \( Y \). Since \( f \) changes by exactly one along edges, we must have that

\[
f(X) = \{1\} \quad \text{and} \quad f(Y) = \{0\}.
\]

Also, since \( f(v) > t \) and \( f \) changes by exactly one along edges, we must have \( B(v, t) \subseteq A \cup X \), which implies that

\[
|X| \geq d(d-1)^{t-1}.
\]
From $f$ we can define many other homomorphisms: for every $s \in \{-1, 1\}$, we can define

$$h_s(w) = \begin{cases} 
  f(w) - 2 & w \in A, \\
  s_w & w \in X, \\
  f(w) & w \notin A \cup X,
\end{cases}$$

and it follows easily from (1.6) that $h_s$ is indeed a homomorphism. Thus we define

$$T(f) = \{h_s : s \in \{-1, 1\}\}.$$

We now study the few-to-many property of $T$. It is convenient to partition the set of functions $P$ according to the sizes of $A$ and $X$. We let

$$P_{x,\beta} := \{f \in P : |A(f)| = x \text{ and } |X(f)| = \beta\}.$$ 

We claim that $T$ is few-to-many from $P_{x,\beta}$ to $Q$. On the one hand,

$$|T(f)| = 2^{|X(f)|} = 2^\beta$$

(1.8)

for all $f \in P_{x,\beta}$. On the other hand, to recover $f$ from $h_s$ we just need to describe the set $A(f)$ (since given $A$ we know $X$, $f$ on $A$ is a translate of $h_s$, $f$ on $X$ is 1, and $f$ on the complement of $A \cup X$ equals $h_s$). Since $A$ is connected in $G \leq 2$, the number of possible sets $A$ of a given size $x$ is at most $d^{4x}$ (see Lemma 2.1 below). Thus, for every $h$ in $Q$,

$$|\{f \in P_{x,\beta} : h \in T(f)\}| \leq d^{4x}. \quad (1.9)$$

Now, we use the fact that subsets of (the vertex set of) the tree which do not contain leaves expand very well (see Claim 4.1 below). Since $A(f)$ does not contain leaves by its definition and the fact that $f$ equals zero on the leaves, we have

$$(d - 2)|A(f)| < |X(f)|$$

(1.10)

for all $f \in P$. Putting together (1.8), (1.9) and (1.10) and applying Lemma 1.8 gives

$$\frac{|P_{x,\beta}|}{|Q|} \leq \frac{d^{4x}}{2^\beta} \leq \frac{d^{4\beta/(d-2)}}{2^\beta} \leq 2^{-\beta/2}$$

for large $d$. By summing over all possibilities for $x$ and $\beta$ and using (1.7), we may conclude that

$$\mathbb{P}(f(v) > t) = \frac{|P|}{|Q|} \leq \exp\left(-\frac{1}{5}d(d-1)^{t-1}\right),$$

for all large $d$, as we wanted to prove.

### 1.5. Readers’ guide

Section 2 gives some preliminary results on graphs. Sections 3 and 4 contain our investigation of the behaviour of Lipschitz functions on good expanders and on trees, respectfully. In Section 5, we study the behaviour of homomorphisms on good bi-expanders. Finally, Section 6 consists of a short summary and suggestions for future research.
2. Preliminaries

We start by fixing some notation. In the following, $G$ is an arbitrary graph. We recall that the graph $G^{\leq 2}$ is defined as follows: it has the same vertex set as $G$ and $\{u,v\} \in E(G^{\leq 2})$ if and only if $dist_G(u,v) \leq 2$, where $dist_G$ is the graph distance in $G$, that is, the length of a shortest path connecting a pair of vertices in $G$. For a set $A$ of vertices of $G$, define the neighbourhood of $A$ by

$$N(A) = \{b \in V(G) : \text{there exists an } a \in A \text{ so that } \{a,b\} \in E(G)\}.$$ 

Define the outer boundary of $A$ as

$$\partial(A) = N(A) \setminus A.$$ 

The 2-outer boundary of $A$ is

$$\partial^2(A) = N(N(A)) \setminus (A \cup N(A)).$$

Denote balls in (the graph metric on) $G$ by

$$B(v,t) = \{w \in V(G) : dist_G(v,w) \leq t\}.$$ 

Finally, for a finite set $X$, denote by $\mathcal{P}(X)$ the family of all subsets of $X$.

In the remainder of this section, we list and prove some preliminary, simple facts.

2.1. Counting connected subgraphs

The following lemma uses depth-first search to bound the number of connected sets in a graph.

**Lemma 2.1.** Let $a$ be an integer, let $H$ be an arbitrary graph with maximum degree $\Delta$ and let $v \in V(H)$. The number of connected sets $A \subseteq V(H)$ such that $v \in A$ and $|A| = a$ does not exceed $\Delta^2a - 2$.

**Proof.** For every such $A$, we fix an arbitrary spanning tree $T_A$ of $A$. We perform a depth-first search on $T_A$, starting and ending at $v$ and passing through every edge exactly twice. Since every spanning tree of $A$ has exactly $a - 1$ edges, the number of possibilities for such a walk (and hence for $A$) is not larger than the number of walks of length $2a - 2$ in $H$ that start at $v$. \hfill {$\square$}

2.2. Expanders

In this section we establish some basic properties of $\lambda$-expanders. Some of them have already been mentioned in our discussion in Section 1.1. Our proof of Lemma 1.1 relies on the fact that every pair of sufficiently large vertex sets in a $\lambda$-expander is connected by an edge, which follows directly from Definition II.

**Proposition 2.2 (connectivity).** Let $G$ be a $d$-regular $n$-vertex $\lambda$-expander. Then, for every two sets $A, B \subseteq V(G)$ satisfying $\min\{|A|, |B|\} > \frac{\lambda n}{d}$, there is an edge of $G$ joining $A$ and $B$.

**Proof.** Since $\frac{d}{n}|A||B| > \lambda \sqrt{|A||B|}$, it follows from Definition II that $e(A,B) > 0$. \hfill {$\square$}
We can now define the phase of a function and prove its properties.

**Proof of Lemma 1.1.** Let $G$ be a $d$-regular $n$-vertex $\lambda$-expander and let $v_0$ be an arbitrary vertex of $G$. Fix an $f \in \text{Lip}_{v_0}(G; M)$ and let $k$ be the smallest integer such that

$$\left| \{ v \in V(G) : f(v) \leq k \} \right| > \frac{\lambda n}{d}. \quad (2.1)$$

Since $f$ is $M$-Lipschitz, there are no edges in $G$ between the sets $\{ v : f(v) \leq k \}$ and $\{ v : f(v) > k + M \}$. It follows from Proposition 2.2 and (2.1) that the latter set has at most $\frac{\lambda n}{d}$ elements. Hence, by minimality of $k$,

$$\left| \{ v : f(v) \notin \{ k, \ldots, k + M \} \} \right| = \left| \{ v : f(v) < k \} \right| + \left| \{ v : f(v) > k + M \} \right| \leq 2 \frac{\lambda n}{d}.$$

This shows that the set of integers $k$ satisfying (1.3) is non-empty for all $f$. We now describe a way to define phase($f$) so that phase($-f$) = $-\text{phase}(f)$ for $f$ which is not identically zero. Fix some total order on Lip$_{v_0}(G; M)$ (e.g., the lexicographic order). For every $f \in \text{Lip}_{v_0}(G; M)$ which is not identically zero, let $b(f)$ be the larger in this total order between $f$ and $-f$, and let $s(f)$ be the smaller of the two. Define phase($b(f)$) as the interval $\{ k, \ldots, k + M \}$ for the minimal $k$ satisfying (1.3) for $b(f)$ and define phase($s(f)$) = $-\text{phase}(b(f))$.

The following standard proposition shows vertex expansion in expanders.

**Proposition 2.3 (expansion).** Let $G$ be a $d$-regular $n$-vertex $\lambda$-expander. Then, for every $A \subseteq V(G)$,

$$|N(A)| \geq \min \left\{ \frac{n}{2}, \frac{d^2}{4\lambda^2} |A| \right\}.$$

**Proof.** Let $A \subseteq V(G)$. We may assume that $A \neq \emptyset$ and $|N(A)| < n/2$. Since there are no edges between $A$ and the complement of $N(A)$, then by Definition II,

$$d|A| = e(A, N(A)) \leq \frac{d}{n} |A||N(A)| + \lambda \sqrt{|A||N(A)|} \leq \frac{d|A|}{2} + \lambda \sqrt{|A||N(A)|}. \quad \square$$

Proposition 2.3 has the following immediate corollary.

**Corollary 2.4 (large boundary).** Let $G$ be a $d$-regular $n$-vertex $\lambda$-expander. Then for every $A \subseteq V(G)$ with $|A| \leq \frac{n}{4}$,

$$|\partial(A)| \geq \min \left\{ \frac{n}{4}, \left( \frac{d^2}{4\lambda^2} - 1 \right) |A| \right\}.$$

The following proposition gives an estimate on the growth of balls in expanders.
Proposition 2.5 (volume growth). Let $G$ be a $d$-regular $n$-vertex $\lambda$-expander. Then for every non-negative integer $t$ and every $v \in V(G)$,

$$|B(v,t)| \geq \min \left\{ \frac{n}{2}, \left( \frac{d}{2\lambda} \right)^{2t} \right\}.$$  \hspace{1cm} (2.2)

Proof. Fix $v$ and prove (2.2) by induction on $t$. The bound trivially holds for $t = 0$. Assume that (2.2) holds for $t \geq 0$. Since $B(v,t+1) \supseteq N(B(v,t))$, Proposition 2.3 implies that

$$|B(v,t+1)| \geq \min \left\{ \frac{n}{2}, \frac{d^2}{4\lambda^2} |B(v,t)| \right\}.$$ 

Proposition 2.5 implies the following bound on the diameter of expanders.

Corollary 2.6 (diameter). Let $G$ be a $d$-regular $n$-vertex $\lambda$-expander. If $\lambda < d/2$, then

$$\text{diam}(G) \leq \left( \log \frac{d}{2\lambda} \right)^{-1} \cdot \log n.$$ 

Proof. Let $v$ and $w$ be two arbitrary vertices of $G$ and let $t$ be the minimum integer such that both $B(v,t)$ and $B(w,t)$ contain more than $\lambda n/d$ vertices. Since $\lambda n/d < n/2$, by Proposition 2.5,

$$t \leq \left( 2 \log \frac{d}{2\lambda} \right)^{-1} \cdot \log \left( \frac{\lambda n}{d} \right) \leq \left( 2 \log \frac{d}{2\lambda} \right)^{-1} \cdot \log n - \frac{1}{2}.$$ 

By Proposition 2.2, $B(v,t)$ and $B(w,t)$ are joined by an edge and hence $\text{dist}_G(v,w) \leq 2t + 1$. \hfill \Box

2.3. Bi-expanders

In this section, we present some basic properties of $\lambda$-bi-expanders. As most of these properties are natural bipartite analogues of the statements presented in Section 2.2 (and given the obvious similarity between Definitions II and V), we leave most of the proofs as an exercise for the reader.

Proposition 2.7 (connectivity). Let $G$ be a $d$-regular $2n$-vertex $\lambda$-bi-expander and let $V_0$ and $V_1$ be the two colour classes of $G$. Then, for every two sets $A \subseteq V_0$ and $B \subseteq V_1$ satisfying

$$\min\{|A|,|B|\} > \frac{\lambda n}{d},$$

there is an edge of $G$ joining $A$ and $B$.

We can now define the phase of a homomorphism function on a $\lambda$-bi-expander and prove its properties.

Proof of Lemma 1.4. Let $G$ be a $d$-regular $2n$-vertex $\lambda$-bi-expander, let $V_0$ and $V_1$ be the two colour classes of $G$, and let $v_0$ be an arbitrary vertex of $G$. Fix $f \in \text{Hom}_{v_0}(G)$, let
$k$ be the smallest integer such that
\[ |\{v \in V_{i-1} : f(v) < k\}| > \frac{\lambda n}{d} \] (2.3)
for some $i \in \{0,1\}$ and fix (the unique) such $i$. Since $f$ is a homomorphism, there are no edges in $G$ between the sets $\{v \in V_{i-1} : f(v) < k\}$ and $\{v \in V_i : f(v) > k\}$. It follows from Proposition 2.7 and (2.3) that the latter set has at most $\frac{\lambda n}{d}$ elements. Since all values taken by $f$ on $V_i$ have the same parity, it follows from the minimality of $k$ that $\{v \in V_i : f(v) < k\} = \{v \in V_i : f(v) < k - 1\}$ and hence
\[ |\{v \in V_i : f(v) \neq k\}| = |\{v \in V_i : f(v) < k - 1\}| + |\{v \in V_i : f(v) > k\}| \leq \frac{2\lambda n}{d}. \]

Finally, suppose that $\lambda < d/3$. Let $k = \text{phase}(f)$ and let $i = i^*$ be such that (1.3) holds. Since $f$ is a homomorphism, there are no edges in $G$ between the sets $\{v \in V_i : f(v) = k\}$ and $\{v \in V_{i-1} : |f(v) - k| \geq 2\}$. By the definition of phase and the assumption on $\lambda$, it follows that the former set has more than $\frac{\lambda n}{d}$ elements and hence, by Proposition 2.7, the latter set has at most $\frac{\lambda n}{d}$ elements. Therefore,
\[ |\{v \in V(G) : |f(v) - k| \geq 2\}| = |\{v \in V_i : f(v) \neq k\}| + |\{v \in V_{i-1} : |f(v) - k| \geq 2\}| \leq \frac{3\lambda n}{d}. \]

Proposition 2.8 (expansion). Let $G$ be a $d$-regular $2n$-vertex $\lambda$-bi-expander. Then for every $A \subseteq V(G)$,
\[ |N(A)| \geq \min\left\{ \frac{n}{2}, \frac{d^2}{4\lambda^2} |A| \right\}. \]

Corollary 2.9 (large boundary). Let $G$ be a $d$-regular $2n$-vertex $\lambda$-bi-expander. Then for every $A \subseteq V(G)$ with $|A| \leq \frac{n}{4}$,
\[ |\partial(A)| \geq \min\left\{ \frac{n}{4}, \left( \frac{d^2}{4\lambda^2} - 1 \right) |A| \right\}. \]

Proposition 2.10 (volume growth). Let $G$ be a $d$-regular $2n$-vertex $\lambda$-bi-expander. Then for every non-negative integer $t$ and every $v \in V(G)$,
\[ |B(v,t)| \geq \min\left\{ \frac{n}{2}, \left( \frac{d}{2\lambda} \right)^{2t} \right\}. \] (2.4)

Corollary 2.11 (diameter). Let $G$ be a $d$-regular $2n$-vertex $\lambda$-bi-expander. If $\lambda \leq d/8$, then
\[ \text{diam}(G) \leq \left( \log \frac{d}{2\lambda} \right)^{-1} \cdot \log n + 1. \]
3. Lipschitz functions

Assume that $G$ is a connected $d$-regular $n$-vertex $\lambda$-expander that is $M$-good, which roughly means that $\lambda \ll d/(M \log d)$. Let $v$ and $v_0$ be two (not necessarily distinct) vertices of $G$ and let $t$ be a positive integer. Recall the definition of phase from Lemma 1.1. We will estimate the probability that a uniformly chosen random function $f \in \text{R Lip}_{v_0}(G; M)$ is in the event

$$\Omega = \{ f \in \text{R Lip}_{v_0}(G; M) : \text{dist}(f(v), \text{phase}(f)) > (t-1)M \}.$$

We first briefly describe our strategy. Our proof of Theorem 1.2 is divided into two (independent) parts.

In the first part, given in Sections 3.1 and 3.2, we construct a map $T : \Omega \to \mathcal{P}(\text{R Lip}_{v_0}(G; M))$ such that the set $T(f)$ is large for every $f \in \Omega$. Moreover, for every $g$, we bound the size of the set $\{ f \in \Omega : g \in T(f) \}$. This is crucial in estimating the probability of $\Omega$ using Lemma 1.8. In this part of the proof, we do not use the assumption that $G$ is a $\lambda$-expander.

In the second part of the proof, Section 3.3, we derive the claimed bound on $\mathbb{P}(\Omega)$ using the properties of the transformation $T$ and the underlying graph $G$. Here, we strongly use the assumption that $G$ is a good expander.

In fact, we partition $\Omega$ into two parts and argue as above on each part:

$$\Omega^+ = \{ f \in \Omega : f(v) > k \} \quad \text{and} \quad \Omega^- = \{ f \in \Omega : f(v) < k \}.$$

Since $\text{phase}(-f) = -\text{phase}(f)$ by the definition of the phase, the map $f \mapsto -f$ is a bijection between $\Omega^+$ and $\Omega^-$. Thus,

$$\mathbb{P}(\Omega) \leq \mathbb{P}(\Omega^+) + \mathbb{P}(\Omega^-) \leq 2\mathbb{P}(\Omega^+). \tag{3.1}$$

Hence, from now on we can focus our attention on the event $\Omega^+$.

3.1. Constructing the transformation $T$

Let $f \in \Omega^+$ and observe that $f$ is not the zero function. In other words, let $f \in \text{R Lip}_{v_0}(G; M)$ satisfy $f(v) > k + tM$, where $k$ is the unique integer such that $\{k, \ldots, k + M\} = \text{phase}(f)$. To make our argument more general, for the remainder of this section and in Section 3.2, let us disregard the fact that $G$ is an $M$-good expander and the precise definition of $k$.

Let us only assume that $G$ is an arbitrary finite connected graph with two (not necessarily distinct) fixed vertices $v_0$ and $v$, that an arbitrary function $k : \text{R Lip}_{v_0}(G; M) \to \mathbb{Z}$ is given and that $\Omega^+$ is a subset of all $f \in \text{R Lip}_{v_0}(G; M)$ satisfying $f(v) > k + M$ for $k = k(f)$. Our only requirement on the function $k$ is that for each $f \in \text{R Lip}_{v_0}(G; M)$ there exists some vertex $w$ on which $f(w) \leq k + M$ for $k = k(f)$. Let

$$A(f) \text{ be the connected component of the vertex } v \text{ in the subgraph of } G^{\leq 2} \text{ induced by the set of vertices } \{ w : f(w) > k + M \}.$$ 

Our requirement on $k$ implies that $A(f)$ does not contain all vertices of $G$. Let us further partition the event $\Omega^+$. Let $C_v$ be the family of all sets of vertices $A \subseteq V(G)$ such that
v ∈ A and A is connected in \( G^{\leq 2} \). For every \( A ∈ C_v \), let

\[
Ω_A^+ = \{ f ∈ Ω^+ : A(f) = A \}.
\]

**Claim 3.1.** Let \( A ∈ C_v \) and \( f ∈ Ω_A^+ \). Set \( X = \partial(A) \). Then, the following properties hold:

1. \( \min f(A) > k + M \),
2. \( f(X) \subseteq \{ k + 1, k + 2, \ldots, k + M \} \) and
3. \( \max f(\partial^2(A)) ≤ k + M \).

**Proof.** (1) Follows since \( A \) is a subset of the set of vertices \( w \) such that \( f(w) > k + M \).

(2) Since \( A \) is defined as a connected component, if \( w \) in \( X \) satisfies \( f(w) > k + M \), then \( w \) is in \( A \) as well, which is a contradiction (\( X ∩ A = \emptyset \)). Since \( f \) is \( M \)-Lipschitz and \( \min f(A) > k + M \), we have \( \min f(X) ≥ \min f(N(A)) > k \).

(3) Similarly to (2), if \( w ∈ \partial^2(A) \) and satisfies \( f(w) > k + M \), then, since \( A \) is defined as a connected component in \( G^{\leq 2} \), we have \( w ∈ A \), a contradiction. □

For the discussion, fix an \( f ∈ Ω_A^+ \) for some \( A ∈ C_v \) and set \( X = \partial(A) \). Here is a first hint on how to define \( T(f) \). For any \( s ∈ \mathbb{Z}^X \), define \( h_s = h_s(f) \), a map from \( V(G) \) to \( \mathbb{Z} \), by

\[
h_s(w) = \begin{cases} 
  k + M & w ∈ A, \\
  k + s_w & w ∈ X, \\
  f(w) & w ∉ A ∪ X.
\end{cases}
\]

Clearly, not every \( s ∈ \mathbb{Z}^X \) gives rise to an \( M \)-Lipschitz function \( h_s \). Still, it is quite easy to identify a large subset of \( \mathbb{Z}^X \) that does have this property. To this end, for every \( x ∈ X \), let

\[
u_x = u_x(f) = \min \{ f(w) + M - k : w ∈ N(x), w ∉ A ∪ X ∪ \{ M \} \}
\]

and

\[
ℓ_x = ℓ_x(f) = \max \{ f(w) - M - k : w ∈ N(x) ∩ A \}.
\]

The following propositions clarifies the relation between the sequences defined above.

**Proposition 3.2.** Let \( A ∈ C_v \), let \( f ∈ Ω_A^+ \) and let \( X = \partial(A) \). Then, for every \( x ∈ X \),

\[
1 ≤ ℓ_x(f) ≤ f(x) - k ≤ u_x(f) ≤ M.
\]

(3.2)

The proposition shows that \( u_x \) is an upper bound on the value of \( f(x) - k \) and \( ℓ_x \) is a lower bound.

**Proof.** Fix \( x ∈ X \). There are four inequalities to prove. First, let \( w_x \) be an element of \( N(x) ∩ A \) such that \( ℓ_x(f) = f(w_x) - M - k \). Since \( w_x ∈ A \), it follows from (1) in Claim 3.1
Claim 3.3. For every \( f \) in \( \Omega^+ \) and for every \( s \in S(f) \), the map \( h_s = h_s(f) \) is an \( M \)-Lipschitz function (but there is no guarantee that \( h_s(v_0) = 0 \)).

Proof. Fix \( s \in S \) and let \( \{w_1, w_2\} \) be an edge of \( G \). It suffices to show that \( |h_s(w_1) - h_s(w_2)| \leq M \). In order to prove this inequality, we consider several cases depending on the locations of the vertices \( w_1 \) and \( w_2 \).

Case 1. If \( w_1, w_2 \in A \), then \( |h_s(w_1) - h_s(w_2)| = 0 \).

Case 2. If \( w_1, w_2 \in X \), then \( |h_s(w_1) - h_s(w_2)| = |s_{w_1} - s_{w_2}| \leq \max \{u_{w_1}, u_{w_2}\} \leq M \).

Case 3. If \( w_1, w_2 \notin A \cup X \), then \( |h_s(w_1) - h_s(w_2)| = |f(w_1) - f(w_2)| \leq M \).

Case 4. If \( w_1 \in A \) and \( w_2 \in X \), then \( |h_s(w_1) - h_s(w_2)| = |k + M - k - s_{w_2}| \leq M \).

Case 5. If \( w_1 \in X \) and \( w_2 \notin A \cup X \), then, by definition of \( u_{w_1} \),

\[
h_s(w_1) - h_s(w_2) = k + s_{w_1} - f(w_2) \leq k + u_{w_1}(f) - f(w_2) \leq M.
\]

On the other hand, by (3) in Claim 3.1,

\[
h_s(w_1) - h_s(w_2) = k + s_{w_1} - f(w_2) \geq k - k - M = -M.
\]

Since there are no edges between \( A \) and \( (A \cup X)^c \), the proof is now complete. \( \square \)

One might be tempted to suggest \( T(f) = \{h_s : s \in S(f)\} \). Although this is a reasonable guess, one needs to be somewhat careful. As \( v_0 \) is arbitrary, it might happen that \( h_s(v_0) \neq 0 \) for some \( s \in S \). Luckily, the following claim reassures us that this is not a serious issue. For this, we define the ‘shift’ operator \( P_{v_0} \) on functions \( h : V(G) \to \mathbb{Z} \) by

\[
P_{v_0}(h) = h - h(v_0).
\]

Clearly, \( P_{v_0}(h)(v_0) = 0 \), and if \( h \) is \( M \)-Lipschitz, so is \( P_{v_0}(h) \).

Claim 3.4. Let \( f \in \Omega^+ \) and \( S = S(f) \). The shift \( P_{v_0} \) is one-to-one on \( \{h_s : s \in S\} \).

Proof. Let \( s \in S \). Since \( h_s(v) = k + M \) and \( h_s - P_{v_0}(h_s) \) is a constant function,

\[
h_s = P_{v_0}(h_s) + (k + M - P_{v_0}(h_s)(v)).
\]

\( \square \)
Finally, define $T : \Omega^+ \to \mathcal{P}(\text{Lip}_0(G;M))$ by

$$T(f) = \{P_{v_0}(h_s) : s \in S(f)\}, \quad (3.3)$$

### 3.2. Properties of $T$

Before we establish the key properties of $T$, we need to further refine our partition of $\Omega^+$. For every $A \in \mathcal{C}_t$ let

$$S(A) = \{S(f) : f \in \Omega^+_{A,t}\},$$

and for every $S \in S(A)$ let

$$\Omega^+_A S = \{f \in \Omega^+_A : S(f) = S\}.$$

We are now ready to prove a key lemma.

**Lemma 3.5.** Let $A \in \mathcal{C}_t$ and $X = \partial(A)$. For every $S = \prod_{x \in X} \{0, \ldots, u_x\}$ in $S(A)$, the following holds.

1. If $f \in \Omega^+_A S$, then $|T(f)| = |S|$.
2. For every $h \in \text{Lip}_0(G;M)$, we have

$$|\{f \in \Omega^+_A S : h \in T(f)\}| \leq M(2|A| + 1)(2M + 1)^d|S_\gamma|,$$

where

$$S_\gamma = \prod_{x \in X} \{1, \ldots, u_x\}.$$

**Proof.** (1) By (3.3) and Claim 3.4,

$$|T(f)| = |\{P_{v_0}(h_s) : s \in S\}| = |S|.$$

(2) Fix some $h \in \text{Lip}_0(G;M)$. Assume that we are given an integer $k$ and some function $f_0 : A \cup X \to \mathbb{Z}$. We claim that there is at most one $f \in \Omega^+_A S$ such that:

- $h \in T(f)$,
- $k = k(f)$ and
- $f(w) = f_0(w)$ for all $w \in A \cup X$.

It suffices to check that if there exists such an $f$, then we can uniquely reconstruct its values on all vertices $w \notin A \cup X$ using only the fact that $h \in T(f)$ and the data in $h, f_0$ and $k$. Recall that $h = P_{v_0}(h_s)$ for some $s \in S$ and that $f(w) = h_s(w)$ for all $w \notin A \cup X$. Hence, it suffices to reconstruct $h_s(v_0)$. But $h_s(v_0) = h_s(v) - h(v) = k + M - h(v)$.

Therefore, in order to bound the number of possible $f$s with $h \in T(f)$, it suffices to bound the number of pairs $(k, f_0)$ for which there exists an $f$ as above.

First, we bound the number of possibilities for $k$. If $v_0 \notin A \cup X$, then since $h_s(v_0) = f(v_0) = 0$, we have $h(v) = k + M$ and hence $k$ is uniquely determined by $h$. If $v_0 \in A \cup X$, observe that since any $f \in \Omega^+_A S$ satisfies $f(v_0) = 0$ and can change by at most $M$ along each edge of $G$, there exists some $x_0 \in X$ such that $|f(x_0)| \leq M|A|$ for all $f \in \Omega^+_A S$. It
follows from part (2) of Claim 3.1 that \( k \in [-M|A| - M, M|A| - 1] \), and hence there are at most \( M(2|A| + 1) \) options for \( f(x) \) for \( f \in \Omega_{A,S}^+ \).

Now, fix \( k \) and bound the number of possibilities for \( f_0 \). Start by bounding the number of possible values of \( f_0(x) \) on \( X \). Let \( f \in \Omega_{A,S}^+ \) be such that \( h \in T(f) \). Since \( S = S(f) \) by the definition of \( \Omega_{A,S}^+ \), then by Proposition 3.2, \( 1 \leq f(x) - k \leq u_x \) for every \( x \in X \). In other words, \( (f_0(x) - k)_{x \in X} \in S_- \).

Finally, we bound the number of possible values of \( f_0 \) on \( A \). Fix some \( s \in S_- \) and suppose that \( f_0(x) = k + s_x \) for every \( x \in X \). We bound the number of ways we can extend an \( M \)-Lipschitz \( f_0 \) from \( X \) to \( A \cup X \). Let \( A' \subseteq A \) be a connected component of \( A \) in \( G \). Observe that to specify \( f_0 \) on \( A' \) it suffices to fix a spanning tree of \( A' \) in \( G \) and specify the difference of values of \( f_0 \) on the edges of this spanning tree and on a single edge leading from \( A' \) to \( X \). Since the spanning tree has exactly \( |A'| - 1 \) edges, there are at most \((2M + 1)^{|A'|}\) possibilities to extend \( f_0 \) to \( A' \). Multiplying this quantity over all connected components of \( A \) in \( G \), we see that there are at most \((2M + 1)^{|A|}\) ways to extend \( f_0 \) from \( X \) to \( A \). \( \square \)

Lemma 3.5 already allows us to prove an estimate on \(|\Omega_{A,S}^+|\) using Lemma 1.8.

**Corollary 3.6.** For every \( A \in C_v \) and \( S \in S(A) \),

\[
\frac{|\Omega_{A,S}^+|}{|T(\Omega_{A,S}^+)|} \leq M(2|A| + 1)(2M + 1)^{|A|} \left( \frac{M}{M + 1} \right)^{|\partial(A)|}.
\]

**Proof.** The claimed estimate follows directly from Lemmas 1.8 and 3.5 and the fact that

\[
\frac{|S_-|}{|S|} = \prod_{x \in X} \frac{u_x}{u_x + 1} \leq \left( \frac{M}{M + 1} \right)^{|X|}
\]

with \( X = \partial(A) \). \( \square \)

To derive a bound on \(|\Omega_A^+|\), we use the following property of the transformation \( T \).

**Claim 3.7.** For every \( A \in C_v \) and every \( S \neq S' \) in \( S(A) \),

\[ T(\Omega_{A,S}^+) \cap T(\Omega_{A,S'}^+) = \emptyset. \]

**Proof.** Fix some \( A \in C_v \) and let \( X = \partial(A) \). Let \( f \in \Omega_A^+ \) and fix an arbitrary \( h \in T(f) \). To prove the claim, we show that we can reconstruct \((u_x(f))_{x \in X}\) from \( h \). Recall that \( h_x(w) = f(w) \) for every \( w \not\in A \cup X \). Hence, for \( w \not\in A \cup X \),

\[
f(w) + M - k = h_x(w) - h_x(v) + 2M = h(w) - h(v) + 2M,
\]

where the first equality follows from the fact that \( h_x(v) = k + M \), and the second equality follows from the fact that \( h \) is a ‘shift’ of \( h_x \). Therefore, for every \( x \in X \),

\[
\begin{align*}
u_x(f) &= \min\{(f(w) + M - k : w \in N(x), w \not\in A \cup X) \cup \{M\}) \\
&= \min\{(h(w) - h(v) + 2M : w \in N(x), w \not\in A \cup X) \cup \{M\})
\end{align*}
\]

\( \square \)
Concluding, Corollary 3.6 and Claim 3.7 imply the following bound on the probability of $\Omega^+_A$.

**Corollary 3.8.** For every $A \in C_v$,

$$\mathbb{P}(\Omega^+_A) \leq M(2|A| + 1)(2M + 1)^{|A|} \left( \frac{M}{M + 1} \right)^{|\partial(A)|}.$$ 

### 3.3. Bounding $\mathbb{P}(\Omega)$

**Proof of Theorem 1.2.** Let us again assume that $G$ is a connected $d$-regular $n$-vertex $\lambda$-expander that is $M$-good and recall that $\Omega^+$ is the family of $f \in \text{Lip}_v(G; M)$ that satisfy $f(v) > k + tM$, where $k = k(f) = \min \text{phase}(f)$. Fix some $f \in \Omega^+$, recall the definition of $A(f)$ from Section 3.1 and observe that by (1.1) in Lemma 1.1,

$$|A(f)| \leq |\{w : f(w) > k + M\}| \leq |\{w : f(w) \notin \text{phase}(f)\}| \leq \frac{2\lambda n}{d}.$$

Moreover, since $f(v) > k + tM$ and $f$ is $M$-Lipschitz, it follows that $B_G(v, t - 1) \subseteq A(f)$. Let

$$A_v = \left\{ A \in C_v : B(v, t - 1) \subseteq A \text{ and } |A| \leq \frac{2\lambda n}{d} \right\}.$$

Clearly, the sets $\Omega^+_A$ defined in Section 3.1, with $A \in A_v$, form a partition of $\Omega^+$. We are now ready to derive a bound on $\mathbb{P}(\Omega)$. Essentially, the bound follows from Corollary 3.8 by showing that the expansion properties of $G$ imply that $|\partial(A)|$ is much larger than $|A|$ for $A \in A_v$. For an integer $\alpha$, let

$$A_{v,\alpha} = \{ A \in A_v : |A| = \alpha \}.$$

Since every $A \in A_v$ is connected in $G^{\leq 2}$ and the maximum degree of $G^{\leq 2}$ is at most $d^2$, then by Lemma 2.1,

$$|A_{v,\alpha}| \leq d^{4\alpha}. \tag{3.4}$$

Let $\alpha$ be such that $A_{v,\alpha}$ is non-empty. In particular, $|B(v, t - 1)| \leq \alpha \leq \frac{2\lambda n}{d}$. Let

$$b_\alpha = \min\{|\partial(A)| : A \in A_{v,\alpha}\}.$$

By Corollary 2.4, since $\lambda \leq \frac{d}{8}$ and hence $\alpha \leq \frac{n}{4}$,

$$b_\alpha \geq \min\left\{ \frac{n}{4}, \frac{d^2}{5\lambda^2 \alpha} \right\} \geq \frac{d\alpha}{8\lambda}, \tag{3.5}$$
where the last inequality holds as \( z \leq 2\lambda n/d \). Hence, by Corollary 3.8 and (3.4),

\[
\sum_{A \in \mathcal{A}_{v,z}} \mathbb{P}(\mathcal{A}^+_{\lambda}) \leq d^{k_2} \cdot M(2z + 1)(2M + 1)^z \left( \frac{M}{M + 1} \right)^{b_3} \tag{3.6}
\]

\[
\leq d^{k_2} M^x e^{2x}(3M)^x \exp \left( - \frac{b_3}{M + 1} \right) \leq \exp \left( - \frac{b_3}{M + 1} + 2 \log(9Md^2)z \right)
\]

\[
\leq \exp \left( - b_3 \left( \frac{1}{M + 1} - \frac{16\lambda \log(9Md^2)}{d} \right) \right) \leq \exp \left( - \frac{b_3}{2(M + 1)} \right), \tag{3.7}
\]

where the last two inequalities follow from (3.5) and our assumption that \( \lambda \leq \frac{d}{32(M+1) \log(9Md^2)} \), respectively. Let us adopt the convention that when we write \( \sum_{x} \) we mean a sum over \( \{x : |\mathcal{A}_{v,x}| \neq 0\} \). Since \( \mathcal{A} \cup \partial(\mathcal{A}) = B(v,t) \) for every \( A \in \mathcal{A}_v \), then \( |B(v,t)| \geq x + b_3 \) for every \( x \) as above, and hence (3.1) and (3.6) give

\[
\mathbb{P}(\Omega) \leq 2\mathbb{P}(\Omega^+) = 2 \sum_{A \in \mathcal{A}_v} \mathbb{P}(\Omega^+_{A}) = 2 \sum_{x} \sum_{A \in \mathcal{A}_{v,x}} \mathbb{P}(\Omega^+_{A})
\]

\[
\leq 2 \sum_{x} \exp \left( - \frac{b_3}{2(M + 1)} \right) = 2 \sum_{x} \exp \left( - \frac{b_3 + x - x}{2(M + 1)} \right)
\]

\[
\leq 2 \exp \left( - \frac{|B(v,t)|}{5(M + 1)} \right) \sum_{x} \exp \left( - \frac{3b_3}{10(M + 1)} + \frac{x}{5(M + 1)} \right).
\]

Thus we need only estimate the last sum. Applying (3.5) and our assumption that

\[
\lambda \leq \frac{d}{32(M+1) \log(9Md^2)},
\]

we have

\[
\sum_{x} \exp \left( - \frac{3b_3}{10(M + 1)} + \frac{x}{5(M + 1)} \right) \leq \sum_{x} \exp \left( - \left( \frac{3d}{80\lambda(M + 1)} - \frac{1}{5(M + 1)} \right) x \right)
\]

\[
\leq \sum_{x} \exp \left( - \left( \frac{96 \log(9Md^2)}{80} - \frac{1}{10} \right) x \right)
\]

\[
\leq \sum_{x} \exp(-2x) \leq \frac{1}{2},
\]

as required. \( \square \)

### 4. Lipschitz functions on trees

In this section we prove Theorem 1.7. Recall the definitions of \( T_h^d, v_r, V_L \) and \( \text{Lip}_{V_L}(T_h^d; M) \) from Section 1.3. Fix integers \( M \geq 1, d \geq 40(M+1) \log(M+1), t \geq 1 \) and a non-leaf vertex \( v \) of \( T_h^d \) (since the theorem is trivial for leaf vertices). We will estimate the probability that a uniformly chosen random function \( f \in \text{Lip}_{V_L}(T_h^d; M) \) is in the event

\[
\Omega = \{ f \in \text{Lip}_{V_L}(T_h^d; M) : |f(v)| > tM \}.
\]
By symmetry, \( \mathbb{P}(\Omega) = 2\mathbb{P}(\Omega^+) \), where
\[
\Omega^+ = \{ f \in \text{Lip}_V(T_d^h; M) : f(v) > tM \}.
\]
Hence it suffices to bound the probability of \( \Omega^+ \). It is convenient to introduce an auxiliary graph \( \tilde{T}_d^h \) by taking the graph \( T_d^h \) and gluing the set of leaves \( V_L \) to one new vertex \( v_0 \). It is clear from the definition that the probability distribution of \( f(v) \) is the same when \( f \in \text{R Lip}_V(T_d^h; M) \) and when \( f \in \text{R Lip}_{v_0}(\tilde{T}_d^h; M) \). Hence we may focus on bounding the probability of the event
\[
\tilde{\Omega}^+ = \{ f \in \text{Lip}_{v_0}(\tilde{T}_d^h; M) : f(v_0) > (t - 1)M \}.
\]
We may now apply the results of Sections 3.1 and 3.2 to the graph \( \tilde{T}_d^h \) with the function \( k(f) \equiv 0 \) defined on \( \text{Lip}_{v_0}(\tilde{T}_d^h) \). In particular, defining \( C_v \) and
\[
\tilde{\Omega}^+_A = \{ f \in \tilde{\Omega}^+ : A(f) = A \}
\]
as in Section 3.1, we deduce from Corollary 3.8 that for every \( A \in C_v \),
\[
\mathbb{P}(\tilde{\Omega}^+_A) \leq M(2|A| + 1)(2M + 1)^{|\partial(A)|} \left( \frac{M}{M + 1} \right)^{|\partial(A)|} \quad (4.1)
\]
As in Section 3.3, we again have that for every \( f \in \tilde{\Omega}^+ \), \( B_G(v, t - 1) \subseteq A(f) \). In addition, denoting \( A = A(f) \), we have \( (A \cup \partial(A)) \cap V_L = \emptyset \) by Claim 3.1, since \( f(V_L) = \{0\} \). Thus, letting
\[
\mathcal{A}_v = \{ A \in C_v : B_v \subseteq A \text{ and } (A \cup \partial(A)) \cap V_L = \emptyset \},
\]
the sets \( \tilde{\Omega}^+_A \) with \( A \in \mathcal{A}_v \) form a partition of \( \tilde{\Omega}^+ \). It remains to use (4.1) to bound the probability of \( \tilde{\Omega}^+ \). For this we will need that subsets of the tree which do not contain leaves have large vertex expansion.

**Claim 4.1.** If \( A \) is a non-empty subset of vertices of \( T_d^h \) which does not contain any leaves, then
\[
|\partial(A)| > (d - 2)|A|.
\]

**Proof.** By induction on \( |A| \). If \( |A| = 1 \), then \( |\partial(A)| = d \). If \( |A| > 1 \), let \( w \) be a vertex in \( A \) which is farthest from the root and let \( A' = A \setminus \{w\} \). By the induction hypothesis and our choice of \( w \), \( |\partial(A)| \geq |\partial(A')| - 1 + (d - 1) > (d - 2)|A'| + d - 2 = (d - 2)|A| \).

The above claim also applies to sets \( A \in \mathcal{A}_v \) since subsets \( A \) of vertices of \( T_d^h \) which satisfy \( (A \cup \partial(A)) \cap V_L = \emptyset \) have the same boundary \( \partial(A) \) in both \( T_d^h \) and \( \tilde{T}_d^h \). We continue to define \( \mathcal{A}_{v,\alpha} \) and \( b_\alpha \) exactly as in Section 3.3, and note that even though \( v_0 \) has very high degree in \( \tilde{T}_d^h \), it is still true that \( |\mathcal{A}_{v,\alpha}| \leq d^{b_\alpha} \) as in (3.4).
Combining (4.1), the bound on $|\mathcal{A}_{v,\alpha}|$ and the above claim, and using a similar calculation to (3.6), we have

$$
\sum_{A \in \mathcal{A}_{v,\alpha}} \mathbb{P}(\Omega_A^+) \leq d^{4z} \cdot M(2z + 1)(2M + 1)^z \left( \frac{M}{M+1} \right)^{b_z} \tag{4.2}
$$

$$
\leq \exp \left( -\frac{b_z}{M+1} + 2 \log(9Md^2)z \right) \tag{4.3}
$$

$$
\leq \exp \left( -b_z \left( \frac{1}{M+1} - \frac{2 \log(9Md^2)}{d-2} \right) \right) \leq \exp \left( -\frac{b_z}{2(M+1)} \right), \tag{4.4}
$$

where the last inequality follows by noting that our assumption that $d \geq 40(M+1) \log(M+1)$ implies

$$
4(M+1) \log(9Md^2) + 2 \leq 5(M+1) \log(9Md^2) = 5(M+1) \log(9M) + 10(M+1) \log d
$$

$$
\leq 5(M+1) \log(9M) + \frac{d}{2} \leq 20(M+1) \log(M+1) + \frac{d}{2} \leq d.
$$

We continue exactly as in Section 3.3, using that $b_z > (d-2)z$ and our assumption that $d \geq 40(M+1) \log(M+1)$, and obtain

$$
\mathbb{P}(\Omega) \leq \exp \left( -\frac{|B_{\gamma^d}(v, t)|}{5(M+1)} \right),
$$

where $B_{\gamma^d}(v, t)$ denotes the graph ball of radius $t$ around $v$ in $\tilde{T}^d_h$. It remains to note that if $\text{dist}_{\gamma^d}(v, V_L) \leq t$, then $\mathbb{P}(\Omega)$ is trivially zero since $f$ is $M$-Lipschitz, so that we can replace $B_{\gamma^d}(v, t)$ by $B_{\gamma^d}(v, t)$. Finally we note that when $\text{dist}_{\gamma^d}(v, V_L) \geq t$, we have $|B_{\gamma^d}(v, t)| \geq d(d-1)^{t-1}$.

5. Homomorphisms

Assume that $G$ is a $d$-regular $2n$-vertex $\lambda$-bi-expander that is good ($\lambda \ll d / \log d$). Let $V_0$ and $V_1$ be the two colour classes of $G$ and let $v$ and $v_0$ be two (not necessarily distinct) vertices of $G$. Without loss of generality, we assume that $v_0 \in V_0$. Recall the definition of phase from Lemma 1.4. We estimate the probability that, given an integer $t \geq 2$, a uniformly chosen random function $f \in_R \text{Hom}_{v_0}(G)$ is in the event

$$
\Omega = \{ f \in \text{Hom}_{v_0}(G) : |f(v) - \text{phase}(f)| > t \}.
$$

Our proof of Theorem 1.5 closely follows the proof of Theorem 1.2 given in Section 3. We construct a map $T : \Omega \rightarrow \mathcal{P}(\text{Hom}_{v_0}(G))$ such that the set $T(f)$ is large for every $f \in \Omega$ and the set $\{ f \in \Omega : g \in T(f) \}$ is small for every $g$. We then derive a bound on $\mathbb{P}(\Omega)$ using Lemma 1.8 and some easy counting. We start by splitting $\Omega$ into two parts:

$$
\Omega^+ = \{ f \in \Omega : f(v) > \text{phase}(f) + t \} \quad \text{and} \quad \Omega^- = \{ f \in \Omega : f(v) < \text{phase}(f) - t \}.
$$

By definition of phase, since $G$ is a good expander, for every $f \in \Omega$ there is exactly one $k$ satisfying (1.3) with $i^*$. This implies that the map $f \mapsto -f$ is one-to-one from $\Omega^+$ to $\Omega^-$, and vice versa. We can therefore focus on the event $\Omega^+$. 
5.1. Defining the transformation $T$

For $f \in \Omega^+$, write $k = k(f) = \text{phase}(f)$ and let $A(f)$ be the connected component of the vertex $v$ in the subgraph of $G^{\leq 2}$ induced by the set of vertices $\{w : f(w) > k + 1\}$. We further partition the event $\Omega^+$. Let $C_v$ be the family of all $A \subseteq V(G)$ such that $v \in A$ and $A$ is connected in $G^{\leq 2}$. For every $A \in C_v$, let

$$
\Omega^+_A = \{ f \in \Omega^+ : A(f) = A \}.
$$

Claim 5.1. Let $A \in C_v$ and $f \in \Omega^+_A$. Set $X = \partial(A)$. Then the following properties hold:

1. $\min f(A) > k + 1$,
2. $f(X) = \{k + 1\}$ and
3. $f(\partial^2(A)) = \{k\}$.

The proof is immediate from the definition of $A(f)$ and the definition of homomorphism height functions. Here is a reasonable guess as to how to define $T(f)$ for $f \in \Omega^+$. Write $A = A(f)$ and $X = \partial(A)$, and for any $s \in \{-1, 1\}^X$, define $h_s : V(G) \to \mathbb{Z}$ by

$$
h_s(w) = \begin{cases} 
  f(w) - 2 & w \in A, \\
  k + s_w & w \in X, \\
  f(w) & w \notin A \cup X.
\end{cases}
$$

The following statement is a direct consequence of Claim 5.1. We omit its proof, which is a simple case analysis similar to the proof of Claim 3.3.

Claim 5.2. For every $f \in \Omega^+$ and $s \in \{-1, 1\}^X$, the function $h_s$ is a homomorphism function, i.e., $|h_s(v) - h_s(w)| = 1$ whenever $v$ and $w$ are adjacent in $G$ (but there is no guarantee that $h_s(v_0) = 0$).

Again, we use the ‘shift’ operator $P_{v_0}$ defined by $P_{v_0}(h) = h - h(v_0)$. Recall that $h \in \text{Hom}(G)$ if and only if $P_{v_0}(h) \in \text{Hom}_{v_0}(G)$.

Claim 5.3. For every $f \in \Omega^+$, $P_{v_0}$ is one-to-one on $\{h_s : s \in \{-1, 1\}^X\}$.

Proof. Let $s \in \{-1, 1\}^X$ and let $w \in A \cap N(X)$. Claim 5.1, since $f$ is a homomorphism, tells us that $f(w) = k + 2$. So, $h_s(w) = f(w) - 2 = k$. Since $h_s - P_{v_0}(h_s)$ is a constant function, $h_s = P_{v_0}(h_s) + (k - P_{v_0}(h_s)(w))$. \qed

Finally, define $T : \Omega^+ \to \mathcal{P}(\text{Hom}_{v_0}(G))$ by

$$
T(f) = \{P_{v_0}(h) : s \in \{-1, 1\}^X\}. 
$$

5.2. Properties of $T$

Lemma 5.4. For every $A \in C_v$, the following holds.
(1) If \( f \in \Omega_4^+ \), then \(|T(f)| = 2^{\lceil \sqrt{d} \rceil} \).

(2) For every \( h \in \text{Hom}_0(G) \), we have

\[ |\{ h : h \in T(f) \}| \leq 2. \]

**Proof.** (1) Follows directly from (5.1) and Claim 5.3.

(2) Let \( X = \hat{\mathcal{C}}(A) \) and fix some \( h \in \text{Hom}_0(G) \) and an integer \( k \). We claim that there is at most one \( f \in \Omega_4^+ \) such that \( h \in T(f) \) and \( \text{phase}(f) = k \). First, since \( f \in \Omega_4^+ \), we have \( f(w) = k + 1 \) for all \( w \in X \). Next, recall that \( h = P_{v_0}(h_s) \) for some \( s \in \{-1, 1\}^X \), that \( f(w) = h_s(w) \) for all \( w \not\in A \cup X \), and that \( f(w) = h_s(w) + 2 \) for all \( w \in A \). Hence, it suffices to reconstruct \( h_s \), which amounts to reconstructing \( h_s(v_0) \). For this, fix an arbitrary \( w \in A \cap N(X) \) and observe, as in the proof of Claim 5.3, that \( h_s(v_0) = h_s(w) - h(w) = k - h(w) \).

Therefore, in order to bound the number of possibilities for \( f \), it suffices to bound the number of integers \( k \) for which there exists an \( f \) as above.

Recall that \( h_s(w) = k = h_s(v_0) + h(w) \) for any \( w \in A \cap N(X) \). There are two cases. (i) If \( v_0 \not\in A \cup X \), then \( h_s(v_0) = f(v_0) = 0 \) and so \( k \) is determined by \( h \). (ii) If \( v_0 \in A \cup X \), then \( h_s(v_0) \) can only take values 0 and \(-2 \) and hence \( k \) can take at most 2 values. \( \square \)

Concluding this part of the discussion, Lemmas 1.8 and 5.4 imply the following bound on the probability of \( \Omega_4^+ \).

**Corollary 5.5.** For every \( A \in \mathcal{C}_v \),

\[ \mathbb{P}(\Omega_4^+) \leq 2^{1-\lceil \sqrt{d} \rceil}. \]

### 5.3. Deriving the bound on \( \mathbb{P}(\Omega) \)

**Proof.** [of Theorem 1.5] Fix some \( f \in \Omega^+ \) and let \( k = \text{phase}(f) \). Recall the definition of \( A(f) \) from Section 5.1 and observe that by (1.4) in Lemma 1.4,

\[ |A(f)| \leq |\{ w : f(w) > k + 1 \}| \leq |\{ w : |f(w) - k| \geq 2 \}| \leq \frac{3\lambda n}{d}. \]

Since \( f(v) > k + t \) and \( f \) is a homomorphism, \( B(v, t-1) \subseteq A(f) \). Let

\[ A_v = \left\{ A \in \mathcal{C}_v : B(v, t-1) \subseteq A \text{ and } |A| \leq \frac{3\lambda n}{d} \right\}. \]

Clearly, the sets \( \Omega_4^+ \) defined in Section 5.1, with \( A \in A_v \), form a partition of \( \Omega^+ \).

We are now ready to derive our bound on \( \mathbb{P}(\Omega) \). For an integer \( x \), let

\[ A_{v,x} = \{ A \in A_v : |A| = x \}. \]

Since every \( A \in A_v \) is connected in \( G^{\leq 2} \) and the maximum degree of \( G^{\leq 2} \) is at most \( d^2 \), by Lemma 2.1,

\[ |A_{v,x}| \leq d^{Ax}. \] (5.2)

Let \( x \) be such that \( A_{v,x} \) is non-empty. In particular, \( |B(v, t-1)| \leq x \leq \frac{3\lambda n}{d} \). Let

\[ b_x = \min\{|\hat{\mathcal{C}}(A)| : A \in A_{v,x}\}. \]
By Corollary 2.9 and since $\alpha \leq 3\frac{n}{d}$ and $\lambda \leq \frac{d}{12}$,

$$b_\alpha \geq \min\left\{ \frac{n}{4}, \frac{d^2}{5\lambda^2} \alpha \right\} \geq \frac{d}{12\lambda^2}. \quad (5.3)$$

Hence, by Corollary 5.5, (5.2), (5.3) and since $G$ is a good bi-expander, for every $\alpha$,

$$\sum_{A \in A_{\alpha}} \mathbb{P}(\Omega_A^+) \leq d^{\alpha} \cdot 2^{1-b_\alpha} \leq \exp(4\alpha \log d - (b_\alpha - 1) \log 2) \leq 2 \exp(-b_\alpha/2). \quad (5.4)$$

Finally, a computation analogous to the one given in the end of Section 3.3 shows that (5.4) implies

$$\mathbb{P}(\Omega) \leq 2\mathbb{P}(\Omega^+) \leq \exp\left( -\frac{|B(v, t)|}{3} \right).$$

6. Summary and future work

This work investigates the typical behaviour of Lipschitz functions on expander graphs. The general understanding obtained is that if the graph is a ‘good’ expander (in an exact numerical sense), then a random Lipschitz function on it is unlikely to fluctuate much at any given vertex and is unlikely to have its maximum value larger than $O(\log(\log n))$, where $n$ is the number of vertices in $G$. Similar results are obtained also for homomorphism height functions on bi-expanders and for grounded Lipschitz functions on regular trees.

There are three parameters influencing our results: the regularity $d$ of the graph, the expansion parameter $\lambda$ of the graph and the maximal slope $M$ of our functions. Our results for Lipschitz functions hold under the assumption that

$$\lambda \leq \frac{d}{32(M + 1) \log(9Md^2)}. \quad \text{It is natural to ask how sharp this condition is.}$$

For example, do similar results continue to hold when $\lambda$ is only slightly less than $d$ (e.g., $\lambda = (1 - \epsilon)d$), or when $d$ is very small (e.g., 3-regular expanders)? In another direction, recalling that $\lambda$ cannot be significantly smaller than $\sqrt{d}$, we see that even on the best of expanders, our results are limited to $M < O(\sqrt{d}/\log d)$. Do similar results continue to hold for larger $M$? Here, one may ask the same question for the limiting continuous model (as $M \to \infty$) in which one samples a random continuous Lipschitz function $f : V(G) \to \mathbb{R}$, i.e., a uniform function (in the sense of Lebesgue measure) from the set of real-valued functions satisfying $|f(v) - f(w)| \leq 1$ for adjacent $v$ and $w$, and $f(v_0) = 0$ for some fixed vertex $v_0$.

We remark that our results for grounded Lipschitz functions on trees apply under the weaker assumption that $M < O(d/\log d)$: see Theorem 1.7. In a subsequent work [12], using a different approach than the one used here, we will show that variants of our theorem for $d$-regular trees continue to hold for arbitrary $M$.

Investigating random Lipschitz functions is of interest on many graphs. Can one obtain general necessary and sufficient conditions on the graph for typical Lipschitz functions to be flat? As in the theory of random surfaces, can the behaviour of random Lipschitz
functions be related to the behaviour of the Gaussian free field (see also [2])? As mentioned already in the Introduction, the behaviour of random homomorphism functions is now reasonably understood on the hypercube [7, 10] and on (finite boxes in) the lattice $\mathbb{Z}^d$ with large $d$ [11]. In addition, Engbers and Galvin [5, 4] use entropy methods (first introduced by Kahn [9]) to obtain very general results on graph homomorphisms on hypercube and certain bipartite graphs. In a subsequent work [13], we will show how the methods developed in this paper can be adapted to yield similar results about general graph homomorphisms on expanders. It seems natural to ask what the entropy methods used by Engbers, Galvin and Kahn would yield for the fluctuations of $M$-Lipschitz functions.

Lastly, we mention the tantalizing question of understanding random homomorphism and Lipschitz functions on the two-dimensional lattice $\mathbb{Z}^2$. It is conjectured that flatness no longer holds for this graph but very little is known. Investigation of this model is related to well-known models of statistical physics such as square ice, the 6-vertex model and 3-colourings (the antiferromagnetic 3-state Potts model). See [11] for a discussion and simulation results.

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