A Tool for Custom Construction of QMC and RQMC Point Sets

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Abstract  We present LatNet Builder, a software tool to find good parameters for lattice rules, polynomial lattice rules, and digital nets in base 2, for quasi-Monte Carlo (QMC) and randomized quasi-Monte Carlo (RQMC) sampling over the $s$-dimensional unit hypercube. The selection criteria are figures of merit that give different weights to different subsets of coordinates. They are upper bounds on the worst-case error (for QMC) or variance (for RQMC) for integrands rescaled to have a norm of at most one in certain Hilbert spaces of functions. Various Hilbert spaces, figures of merit, types of constructions, and search methods are covered by the tool. We provide simple illustrations of what it can do.

1 Introduction

QMC methods approximate an integral of the form

$$
\mu = \int_0^1 \cdots \int_0^1 f(u_1, \ldots, u_s) du_1 \cdots du_s = \int_{(0,1)^s} f(u) du = E[f(U)]
$$

(1)
where \( f : (0, 1)^s \rightarrow \mathbb{R} \), by the average
\[
\bar{\mu}_n = \frac{1}{n} \sum_{i=0}^{n-1} f(u_i)
\] (2)

where \( P_n = \{u_0, \ldots, u_{n-1}\} \subset [0, 1]^s \) is a set of \( n \) deterministic points that cover the unit hypercube more evenly than typical independent random points. That is, the discrepancy between their empirical distribution and the uniform distribution over \([0, 1]^s\) is smaller than for independent random points and converges to 0 faster than \( O(n^{-1/2}) \) when \( n \to \infty \). This discrepancy can be defined in many ways. It usually represents the worst-case integration error for a given class of integrands \( f \). Typically, this class is a reproducing-kernel Hilbert space (RKHS) \( \mathcal{H} \) of functions, such that
\[
|E_n| := |\bar{\mu}_n - \mu| \leq \mathcal{D}(P_n)\mathcal{V}(f)
\] (3)

for all \( f \in \mathcal{H} \), where \( \mathcal{V}(f) \) is the norm of \( f - \mu \) in \( \mathcal{H} \) (we call it the variation of \( f \)) and \( \mathcal{D}(\cdot) \) is the discrepancy measure associated with this Hilbert space \([8, 18, 39]\). For a fixed \( f \in \mathcal{H} \) with \( \mathcal{V}(f) > 0 \), the error bound in (3) converges at the same rate as \( \mathcal{D}(P_n) \). A traditional version of (3), whose derivation does not involve Hilbert spaces, is the classical Koksma-Hlawka inequality, in which \( \mathcal{V}(f) \) is the Hardy-Krause variation and \( \mathcal{D}(P_n) \) is the star discrepancy, which converges as \( \mathcal{O}((\log n)^{1/2}n^{-1}) \) for well-selected point sets \([39]\). Another important choice for \( \mathcal{H} \) is a Sobolev space of functions whose mixed partial derivatives of order up to \( \alpha \) are square-integrable. It is known that for this space, one can construct point sets whose discrepancy converges as \( \mathcal{O}((\log n)^{s-1/2}n^{-\alpha}) \), and that this is the best possible rate \([4, 8, 14, 15, 16, 17]\).

The main classes of QMC methods are lattice points and digital nets.

For RQMC, the \( n \) QMC points are randomized to provide a set of random points \( \{U_0, \ldots, U_{n-1}\} \subset (0, 1)^s \) for which (i) each \( U_i \) individually has the uniform distribution over \([0, 1]^s\), and (ii) the points keep their highly-uniform distribution collectively. Randomizations that provably preserve the low discrepancy generally depend on the type of QMC construction: some are used for lattice points and others for digital nets. In some cases, the randomization may even improve the convergence rate of the mean square discrepancy. The RQMC estimator
\[
\hat{\mu}_{n,\text{rqmc}} = \frac{1}{n} \sum_{i=0}^{n-1} f(U_i),
\] (4)

which is now random, is unbiased for \( \mu \) and one wishes to minimize its variance. For more details on RQMC, see for example \([21, 26, 29, 30, 37, 44, 45, 46]\).

The aim of this paper is to introduce \textit{LatNet Builder}, a software tool designed to construct specific point sets for QMC and RQMC, in any number of dimensions, for an arbitrary number of points, arbitrary weights on the subsets of coordinates, arbitrary smoothness of the integrands, a variety of construction and randomization methods, and several choices of discrepancies. The point sets can also be extensible in the number of points and number of dimensions. By “constructing the points” here
we mean defining the set \( P_n \) by selecting the parameter values for a general structure, by trying to minimize a figure of merit (FOM) that may represent a discrepancy \( \mathcal{D}(P_n) \) or be an upper bound on it. Once this is done, other software can be used to randomize and generate the points for their utilization in applications; see [27, 42] for example. LatNet Builder is available in open source at https://github.com/umontreal-simul/latnetbuilder. It is a descendant of Lattice Builder [32], whose scope was limited to ordinary lattice rules. Another related tool is Nuyens’ fast CBC constructions [40].

The rest of this paper is organized as follows. In Section 2, we recall the types of QMC point sets covered by our software, namely ordinary lattice points, polynomial lattice points, digital nets, and their higher-order and interlaced versions, as well as the main randomization methods to turn these point sets into RQMC points. In Section 3, we give the general form of weighted RKHS used in this paper and the corresponding generalized Koksma-Hlawka inequality. We also recall the most common types of weights. In Section 4, we give examples of discrepancies that are used as FOMs to select the parameters of point set constructions. In Section 5, we summarize the search methods implemented in our software. In Section 6, we compare FOM values obtained by various point set constructions and search methods. We also compare RQMC variance for simple integrands \( f \). Section 7 gives a conclusion.

2 Point Set Constructions and Randomizations

LatNet Builder handles ordinary rank-1 lattice points as well as digital nets, which include polynomial lattice rules and high-order and interlaced constructions.

For a rank-1 lattice rule, the point set is

\[
P_n = \{ u_i = iv_1 \mod 1, \; i = 0, \ldots, n-1 \}
\]

where \( n v_1 = a = (a_1, \ldots, a_s) \in \mathbb{Z}_n^s = \{0, \ldots, n-1\}^s \). It is a Korobov rule if \( a = (1, a, a^2 \mod n, \ldots, a^{s-1} \mod n) \) for some integer \( a \in \mathbb{Z}_n \). The parameter to select here is the vector \( a \), for any given \( n \). The usual way to turn a lattice rule into an RQMC point set is by a random shift: generate a single random point \( U \) uniformly in \((0,1)^s\), and add it to each point of \( P_n \), modulo 1, coordinate-wise. This satisfies the RQMC conditions. For more details on lattice rules and their randomly-shifted versions, see [19, 20, 26, 29, 31, 49].

The Digital nets in base 2 handled by LatNet Builder are defined as follows. The number of points is \( n = 2^k \) for some integer \( k \). We select an integer \( w \geq k \) and \( s \) generating matrices \( C_1, \ldots, C_s \) of dimensions \( w \times k \) and of rank \( k \), with elements in \( \mathbb{Z}_2^k \equiv \{0, 1\} \). The points \( u_i, i = 0, \ldots, 2^k - 1 \), are defined as follows: for \( i = a_{i,0} + a_{i,1}2 + \cdots + a_{i,k-1}2^{k-1} \), we take

\[
\begin{pmatrix}
u_{i,1,1} \\
\vdots \\
u_{i,w,1}
\end{pmatrix} = C_j
\begin{pmatrix}
a_{i,0} \\
\vdots \\
a_{i,k-1}
\end{pmatrix} \mod 2,
\]

\[
u_{i,j} = \sum_{\ell=1}^{w} u_{i,j,\ell} 2^{-\ell},
\]
We select a polynomial modulus \( Q \) (which includes PLRs), that can provide the higher-order convergence rate of an interlacing factor \( d \) most to examine each time we select a new coordinate of the generating vector. The most popular digital net constructions are still the Sobol’ points [51], in base \( b = 2 \), with \( k \times k \) generating matrices that are upper triangular and invertible. These matrices are constructed by a specific method, but the bits of the first few columns above the diagonal can be selected arbitrarily, and their choice has an impact on the quality of the net. General-purpose choices have been proposed in [24, 34], e.g., based on the uniformity of two-dimensional projections. LatNet Builder allows one to construct the matrices based on a much more flexible class of criteria.

A polynomial lattice rule (PLR) in base 2 is defined as follows. We denote by \( \mathbb{Z}_2[z] \) the ring of polynomials with coefficients in \( \mathbb{Z}_2 \), by \( \mathbb{L}_2 \) the set of formal series of the form \( \sum_{\ell=\ell_0}^\infty x_\ell z^{-\ell} \) with each \( x_\ell \in \mathbb{Z}_2 \) and \( \ell_0 \in \mathbb{Z} \), and we define \( \varphi : \mathbb{L}_2 \to \mathbb{R} \) by

\[
\varphi \left( \sum_{\ell=\ell_0}^\infty x_\ell z^{-\ell} \right) = \sum_{\ell=\max(\ell_0,1)}^w x_\ell 2^{-\ell}.
\]  

We select a polynomial modulus \( Q = Q(z) \in \mathbb{Z}_2[z] \) of degree \( k \), and a generating vector \( a(z) = (a_1(z), \ldots, a_s(z)) \in \mathbb{Z}_2[z]^s \), whose coordinates are polynomials of degrees less than \( k \) having no common factor with \( Q(z) \). The point set of cardinality \( n = 2^k \) is

\[
P_n = \left\{ \left( \varphi \left( \frac{h(z)a_1(z)}{Q(z)} \right), \ldots, \varphi \left( \frac{h(z)a_s(z)}{Q(z)} \right) \right) : h(z) \in \mathbb{Z}_2[z] \text{ and } \deg(h(z)) < k \right\}.
\]  

Here, we want to optimize the vector \( a(z) \). This point set turns out to be a digital net in base 2 whose generating matrices \( C_j \) contain the first \( w \) digits of the binary expansion of the \( a_j(z)/Q(z) \). These are Hankel matrices: each row is the previous one shifted to the left by one position, with the last entry determined by the recurrence with characteristic polynomial \( Q(z) \), applied to the entries of the previous row. They can have an arbitrary number of rows, but for ordinary PLRs, they are usually truncated to \( w = k \) rows. See [8, 33, 35, 39, 38] for further details on PLRs.

A high-order polynomial lattice rule (HOPLR) of order \( \alpha \) with \( n = 2^k \) points is obtained by constructing an ordinary PLR with polynomial modulus \( \tilde{Q}(z) \) of degree \( \alpha k \) having \( 2^{\alpha k} \) points in \( s \) dimensions, and using only the first \( n = 2^k \) points. See [1, 2, 7]. This type of construction can achieve a higher order of convergence for the error (almost \( O(n^{-\alpha}) \)) than an ordinary PLR for integrands \( f \) in a Sobolev space of smoothness order \( \alpha \) (i.e., when all mixed partial derivatives of up to order \( \alpha \) are square integrable). One drawback is that because of the high degree of \( \tilde{Q} \), the cost of a full CBC construction (see Section 5) is much higher since there are \( 2^{2^k} \) possibilities to examine each time we select a new coordinate of the generating vector.

Dick [3] also proposed an interlacing construction, for digital nets in general (which includes PLRs), that can provide the higher-order convergence rate of almost \( O(n^{-\alpha}) \) for the integration error, for integrands with smoothness order \( \alpha \). For an interlacing factor \( d \in \mathbb{N} \), the method first constructs a digital net in \( sd \) dimen-
sions, with generating matrices \( C_1, \ldots, C_{sd} \). Then the generating matrices of the \( s \)-dimensional interlaced net are \( C^{(d)}_1, \ldots, C^{(d)}_s \), where the rows of \( C^{(d)}_j \) are the first rows of \( C_{(j-1)d+1}, \ldots, C_{jd} \) in this order, then the second rows of these matrices in the same order, and so on.

The simplest way to define a RQMC point set from a digital net in base 2 is to add a digital random shift modulo 2 to all the points. To do this, we generate a single point \( U = (U_1, \ldots, U_s) \) uniformly in \((0, 1)^s\), and perform a bitwise exclusive-or (XOR) between the binary digits of \( U \) and the corresponding digits of each point \( u_i \).

A more involved randomization method for digital nets is the nested uniform scramble (NUS) of Owen [44, 45]. In base 2, for each coordinate, we do the following. With probability 1/2, flip the first bit of all the points. Then, for the points whose first bit is 1, with probability 1/2, flip all the second bits. Do the same for the points whose first bit is 0, independently. Then do this recursively for all the bits. After all flipping is done for the first \( \ell \) bits, partition the points in \( 2^{\ell} \) batches according to the values of their first \( \ell \) bits, and for each batch, flip bit \( \ell + 1 \) of all the points with probability 1/2, independently across the batches. This requires \( (2^{\ell} - 1)s \) random bits to flip the first \( \ell \) bits of all coordinates. One can equivalently do this only for the first \( k \) bits, and generate the other bits randomly and independently across points [37].

A less expensive scramble, which gives less independence than NUS but more than a digital random shift, is a (left) linear matrix scramble (LMS) followed by a digital random shift (LMS+shift) [21, 23, 37, 47]. The LMS replaces \( C_j \) by \( L_j C_j \mod 2 \), where \( L_j \) is a random non-singular lower-triangular \( w \times w \) binary matrix.

Owen [45] proved that under sufficient smoothness conditions on \( f \), the RQMC variance with NUS on digital nets with fixed \( s \) and bounded \( t \) converges as \( O\left(n^{-3}(\log n)^{s-1}\right) \). A variance bound of the same order was shown for LMS+shift in [21, 53]. Note that these results were proved under the assumption that \( w = \infty \).

### 3 Hilbert Spaces and Projection-Dependent Weights

The FOMs used by LatNet Builder are based on generalized (weighted) Koksma-Hlawka inequalities of the form (3) where

\[
\mathcal{V}^{p}(f) = \sum_{\emptyset \neq u \subseteq \{1,2,\ldots,s\}} \gamma_u \mathcal{V}^{p}(f_u)
\]

and

\[
\mathcal{D}^{q}(P_n) = \sum_{\emptyset \neq u \subseteq \{1,2,\ldots,s\}} \gamma_u \mathcal{D}^{q}(P_n),
\]

where \( 1/p + 1/q = 1 \), \( \gamma_u \in \mathbb{R} \) is a weight assigned to the subset \( u \), \( \mathcal{V}(f_u) \) is the variation of \( f_u \), \( \mathcal{D}_u(P_n) \) is the discrepancy of the projection of \( P_n \) over the subset \( u \) of coordinates, and \( f = \sum_{u \subseteq \{1,2,\ldots,s\}} f_u \) is the functional ANOVA decomposition of \( f \) [10, 46]. LatNet Builder allows any \( q \in [1, \infty] \). Taking \( q = \infty \) with \( p = 1 \) means removing the \( q \) and taking the max instead of the sum in (8), while \( p = \infty \) with
$q = 1$ means removing the $p$ and taking the max instead of the sum in (7). The most common choice is $p = q = 2$.

LatNet Builder implements a variety of choices for $D_u(P_n)$, depending on the point set constructions. Some of these measures correspond to the worst-case error in some function space, assuming that the points of $P_n$ are not randomized. Others correspond to the mean-square error (or variance), assuming that the points are randomized in some particular way. This is typically done by defining a RKHS with a kernel that is invariant with respect to the given randomization (i.e., digital shift-invariant, scramble-invariant, etc.), and taking the worst-case error in that space.

The role of the weights is to better recognize the importance of the subsets $u$ for which $f_u$ contributes the most to the error or variance. That is, if $\mathcal{Y}(f_u)$ is unusually large, we want to divide it by a larger weight $\gamma_u$ to control its contribution to $\mathcal{Y}(f)$, but then we have to multiply $D_u(P_n)$ in (8) by the same weight. The final effect is that the FOM will penalize more the discrepancy for that particular projection.

In principle, the weights $\gamma_u$ can be arbitrary. But for large $s$, defining arbitrary individual weights for the $2^s - 1$ projections is impractical, so special forms of weights that are parameterized by much fewer than $2^s - 1$ parameters have been proposed. The most common ones are product weights, for which a weight $\gamma_j$ is assigned to coordinate $j$ for $j = 1, \ldots, s$, and $\gamma_u = \prod_{j \in u} \gamma_j$; order-dependent weights, for which $\gamma_u = \Gamma_{|u|}$, where $\Gamma_1, \ldots, \Gamma_s$ are selected constants and $|u|$ is the cardinality of $u$; and the product-and-order-dependent (POD) weights, which are a combination of the two, with $\gamma_u = \Gamma_{|u|} \prod_{j \in u} \gamma_j$. These are all available in LatNet Builder.

LatNet Builder can construct point sets that are extensible in the number of dimensions and also in the number of points, which means that we can construct point sets that perform well in the first $s$ dimensions for $s = s_{\min}, \ldots, s_{\max}$, and/or if we take the first $n$ points for $n = n_1, n_2, \ldots, n_m$, simultaneously. Typically, one would take $n_j = 2^{k_{\min} + j - 1}$ for $j = 1, \ldots, m$, so $n_m = 2^{k_{\max}} = 2^{k_{\min} + m - 1}$ [21]. The global FOM in this case will be a weighted sum or maximum of the FOMs over the considered dimensions $s$ and/or cardinalities $n_j$. The CBC construction approach described in Section 5 already gives a way to implement the extension in $s$. For the extension in $n$ (or $k$), LatNet Builder implements criteria and heuristic search methods that account for a global FOM.

4 Figures of Merit

Most FOMs considered in LatNet Builder have the general form (8) where typically, when the points have the appropriate special structure of a lattice, polynomial lattice, or digital net, and with an adapted FOM, we have

$$
\mathcal{D}_u^q(P_n) = \frac{1}{n} \sum_{i=0}^{n-1} \prod_{j \in u} \phi(u_{i,j})
$$

(9)
for some function $\phi : [0, 1) \to \mathbb{R}$. With product weights $\gamma_j = \prod_{j \in u} c_j$, the FOM then becomes

$$\mathcal{P}_n(P_n) = -1 + \frac{1}{n} \sum_{j=0}^{n-1} \prod_{j=1}^{s}(1 + \gamma_j \phi(u_{i,j})),$$

which can be computed with $O(ns)$ evaluations of $\phi$.

As an illustration, for a randomly-shifted lattice rule, the variance is:

$$\text{Var}[\hat{\beta}_{n,rqmc}] = \sum_{0 \neq h \in L_n^*} |\hat{f}(h)|^2,$$

where $L_n^* \subset \mathbb{Z}^d$ is the dual lattice [29]. It is also known that for periodic continuous functions having square-integrable mixed partial derivatives up to order $\alpha/2$ for an even integer $\alpha \geq 2$, one has $|\hat{f}(h)|^2 = O((\max(1, h_1) \cdots \max(1, h_s))^{-\alpha})$. This motivates the well-known FOM [31, 39, 49]:

$$\mathcal{P}_\alpha := \sum_{0 \neq h \in L_n^*} (\max(1, h_1) \cdots \max(1, h_s))^{-\alpha}$$

$$= -1 + \frac{1}{n} \sum_{j=0}^{n-1} \prod_{j=1}^{s} \left(1 - \frac{(-4\pi^2)^{\alpha/2}}{\alpha!} B_\alpha(u_{i,j})\right)$$

$$= \frac{1}{n} \sum_{j=0}^{n-1} \sum_{\{1, \ldots, s\} \setminus \{u_{i,j}\}} \left(\frac{-(-4\pi^2)^{\alpha/2}}{\alpha!}\right)^{|u|} \prod_{j \in u} B_\alpha(u_{i,j}),$$

(11)

where $B_{\alpha/2}$ is the Bernoulli polynomial of degree $\alpha/2$ ($B_1(u) = u - 1/2, B_2(u) = u^2 - u + 1/6$, etc.), and the equality in (11) holds only when $\alpha$ is an even integer. Moreover, there are rank-1 lattices point sets $P_n$ for which $\mathcal{P}_\alpha$ converges as $O(n^{-\alpha})$ [9, 48, 49], i.e., as $O(n^{-\alpha + \varepsilon})$ for any $\varepsilon > 0$. Adding projection-dependent weights $\gamma_u$ leads to the weighted $\mathcal{P}_{\gamma,u,\alpha}$, defined by (8) and (9) with $q = 2$,

$$\phi(u_{i,j}) = -(-4\pi^2)^{\alpha/2} B_\alpha(u_{i,j})/\alpha!,$$

and $\mathcal{P}_{\gamma,u}^2(P_n) = \mathcal{P}_{\gamma,u}(P_n)$ is the $\mathcal{P}_\alpha$ for the projection of $P_n$ on the coordinates in $u$.

There is a similar variance expression for digital nets in base 2 with a random digital shift, with the Fourier coefficients $\hat{f}(h)$ replaced the the Walsh coefficients, and the dual lattice replaced by the dual net $\mathcal{D}$ or the dual lattice in the case of PLRs [25, 35]. Thus, FOMs that correspond to variance bounds can be obtained by finding easily computable bounds on the Walsh coefficients. By assuming a rate of decrease of $O(2^{-\alpha|h|})$ of the Walsh coefficients $\hat{f}(h)$ with $h = (h_1, \ldots, h_s) \in \mathbb{N}^s$ and $|h| = |h_1| + \cdots + |h_s|$, and using a RKHS with shift- and scramble-invariant kernel, [53] and [6] obtain a FOM of the form (8) and (9) with

$$\phi(x) = \phi_\alpha(x) = \mu(\alpha) - \mathbb{I}[x > 0] \cdot 2^{(1+|\log_2(x)|)(\alpha-1)} (\mu(\alpha) + 1),$$
where \( \mathbb{I} \) is the indicator function and \( \mu(\alpha) = (1 - 2^{1-\alpha})^{-1} \) for any real number \( \alpha > 1 \). This gives \( \mu(2) = 2, \mu(3) = 4/3, \ldots \). For \( \alpha = 2 \), this gives

\[
\phi_2(x) = 2(1 - \mathbb{I}[x > 0] \cdot 3 \cdot 2^{\log_2(x)}),
\]

which corresponds to the FOM suggested in [35, Section 6.3] for PLRs. In [22, 53], \( \phi(x) \) is written in terms of \( \eta = \alpha - 1 \) instead, but it is exactly equivalent. These papers also show the existence of digital nets for which the FOM converges as \( O(n^{-\alpha}(\log n)^{\epsilon-1}) \) for any \( \alpha > 1 \). This FOM can be seen as a counterpart of \( \mathcal{P}_\alpha \) and we call it \( \tilde{\mathcal{P}}_\alpha \). Its value \( \tilde{\mathcal{P}}_{\alpha,u}(P_n) \) on the projection of \( P_n \) on the coordinates in \( u \) can be used for \( \mathcal{P}^2_{\alpha}(P_n) \), with \( q = 2 \).

Dick and Pillichshammer [8, Chapter 12] consider a RKHS with shift-invariant kernel, which is a weighted Sobolev space of functions whose mixed partial derivatives of order 1 are square-integrable. This gives a FOM of the form (8) and (9) with \( q = 2 \) and

\[
\phi(x) = 1/6 - \mathbb{I}[x > 0] \cdot 2^{\log_2(x)}.
\]

They show that there are digital nets for which this FOM (and therefore the square error) converges almost as \( O(n^{-2}) \). In their Chapter 13, they find that the scramble-invariant version gives the same \( \phi \).

Goda [12] examines interlaced polynomial lattice rules (IPLR), also for a Sobolev space of order \( \alpha \), with an interlacing factor \( d \geq 1 \). He provides two upper bounds on the worst-case error in a deterministic setting. These bounds can be used as FOMs. The first is valid for all positive integer values of \( \alpha \) and \( \delta \), whereas the second holds only for \( d \leq \alpha \), but is tighter when it applies. These two bounds have the form (8) and (9) with \( q = 1, \gamma_u \) replaced by \( \tilde{\gamma}_u \), and

\[
\phi(x, \alpha, d, \ell) = 1 + \prod_{\ell=1}^{d} (1 + \phi_{\alpha,d,\ell}(x_{i,j} - 1, j_{d+\ell})),
\]

where for the first bound, \( \tilde{\gamma}_u = 2^{\alpha(2d-1)/2} \),

\[
\phi_{\alpha,d,\ell}(x) = 1 - 2^{\min(\alpha,d) - 1/\log_2 x} \cdot \left( 2^{\min(\alpha,d) - 1} - 1 \right)
\]

for all \( x \in [0, 1] \), where \( 2^{\log_2 0} = 0 \), while for the second bound, \( \tilde{\gamma}_u = \gamma_u \) and

\[
\phi_{\alpha,d,\ell}(x) = 2^{d-1} \cdot \left( 1 - 2^{(d-1)/\log_2 x} \cdot (2^{d-1} - 1) \right).
\]

One can achieve a convergence rate of almost \( O(n^{-\min(\alpha,d)}) \) for these FOMs (and therefore for the worst-case error).

Goda and Dick [13] proposed another FOM, also for a Sobolev space of order \( \alpha \), for interlaced randomly-scrambled PLRs of high order. They showed that this scheme can achieve the best possible convergence rate of \( O(n^{-\min(\alpha,d)+1} + \delta) \) for the variance. The FOM, which we denote by \( \mathcal{I}_\alpha \), has the same form, but with \( q = 2 \),
\[ \phi(x) = \phi_{\alpha,d}(x) = \frac{1 - 2^{2\min(\alpha,d)[\log_2 x]} (2^{2\min(\alpha,d)+1} - 1)}{1 - 2^{-2\min(\alpha,d)}}, \]

and \( \gamma_u \) replaced by \( \tilde{\gamma}_u = \gamma_u D_{\alpha,d}^{[u]} \) where \( D_{\alpha,d} = 2^{2\max(d - \alpha,0) + (2d - 1)\alpha} \).

Another set of FOMs are obtained from upper bounds on the star discrepancy of \( D^*(P_n) \) or its projections on subsets of coordinates, when \( P_n \) is a digital \((t,k,s)\)-net. One such bound is \( D^*(P_n) \leq 1 - \left(1 - \frac{1}{n}\right)^s + \mathcal{R}_2 \) where

\[ \mathcal{R}_2 = -1 + \frac{1}{n} \sum_{i=0}^{n-1} \prod_{j=1}^{s} \left[ \sum_{k=0}^{n-1} 2^{-|\log_2 k| - 1} \text{wal}_k(u,i,j) \right] \tag{12} \]

\( \text{wal}_k \) is the \( k \)th Walsh function in one dimension, and we assume that the generating matrices \( C_j \) are \( k \times k \). See [8, Theorems 5.34 and 5.36], where a more general version with projection-dependent weights is also given. For PLRs in base \( b = 2 \), this criterion is equal to \( \mathcal{R}_{2,T} \) given in [8, Chapter 10]:

\[ \mathcal{R}_{2,T} = - \sum_{\emptyset \neq u \subseteq \{1,2,\ldots,s\}} \gamma_u + \frac{1}{n} \sum_{i=0}^{n-1} \sum_{\emptyset \neq u \subseteq \{1,2,\ldots,s\}} \gamma_u \prod_{j \in u} \phi_k(u,i,j) \tag{13} \]

where \( \phi_k(u) = -\lfloor \log_2 (u) \rfloor / 2 \) if \( u \geq 2^{-k} \) and \( \phi_k(u) = 1 + k/2 \) otherwise.

A classical upper bound on the star discrepancy is also given by the \( t \)-value of the digital net:

\[ D^*(P_n) \leq 2^t \frac{1}{n} \sum_{j=0}^{s-1} \binom{k-t}{j}. \]

If we use this upper bound for each projection on the subset \( u \) of coordinates, we get the FOM (8) with \( q = 1 \) and

\[ D_u(P_n) = \frac{2^t_u |u|^{-1}}{n} \sum_{j=0}^{t_u - 1} \binom{k-t_u}{j} \]

where \( t_u = t_u(P_n) \) is the \( t \)-value of the projection of \( P_n \) on the coordinates in \( u \). LatNet Builder implements this with arbitrary weights, using algorithms described in [36]. Dick [4] obtains worst-case error bounds that converge at rate almost \( O(n^{-\alpha}) \) for interlaced digital nets, based on the \( t \)-values of the projections.

5 Search Methods

For given construction type, FOM, and weights, finding the best choice of parameters may require to try all possibilities, but their number is usually much too large. LatNet Builder implements the following search methods.
In an exhaustive search, all choices of parameters are tried, so we are guaranteed to find the best one. This is possible only when there are not too many possibilities.

A random search samples uniformly and independently a fixed number \( r \) of parameter choices, and the best one is retained.

In a full component-by-component (CBC) construction, the parameters are selected for one coordinate at a time, by taking into account the choices for the previous coordinates only [5, 8, 50]. The parameters for coordinate \( j \) (e.g., the \( j \)th coordinate of the generating vector in the case of lattices), are selected by minimizing the FOM for the first \( j \) coordinates, in \( j \) dimensions, by examining all possibilities of parameters for this coordinate, without changing the parameter choices for the previous coordinates. This is done for the \( s \) coordinates in succession. This greedy approach can reduce by a huge factor (exponential in the dimension) the total number of cases that are examined in comparison with the exhaustive search. What is very interesting is that for most types of QMC constructions and FOMs implemented in LatNet Builder, the convergence rate for the worst-case error or variance obtained with this restricted approach is provably the same as for the exhaustive search [8, 14].

A random CBC construction can be used when the number of choices for each coordinate is too large: one examines only a fixed number of random choices for each \( j \).

For lattice-type point sets, with certain FOMs, a fast CBC construction can be implemented by using a fast Fourier transform (FFT), so the full CBC construction can be performed much faster [8, 43, 41]. LatNet Builder supports this.

For lattice-type constructions, one can also further restrict the search to Korobov-type generating vectors. The first coordinate is set to 1 and only the second coordinate needs to be selected. This can be done either by an exhaustive search or by just taking a random sample for the second coordinate (random Korobov).

For digital nets, a mixed CBC method is also available: it uses full CBC for the first \( d - 1 \) coordinates and random CBC for the other ones, for a given \( 1 \leq d \leq s \).

### 6 Simple Numerical Illustrations

Here we give a few simple examples of what LatNet Builder can do. The simulation experiments, including the generation and randomization of the points, were done using SSJ [28].

#### 6.1 FOM quantiles for different constructions

One might be interested in estimating the probability distribution of FOM values obtained when selecting parameters at random for a given type of construction, perhaps under some constraints, and as a function of \( n \). Here we estimate this distribution by its empirical counterpart with an independent sample of size 1000
(with replacement), and we report the 0.1, 0.5 and 0.9 quantiles of this empirical distribution, for $n$ going from $2^6$ to $2^{18}$. We do this for PLRs, Sobol’ points, and digital nets with arbitrary invertible and projection-regular generating matrices (random nets), with $\mathcal{P}_2$ taken as the FOM, in $s = 6$ dimensions, with $\gamma_u = 0.7|u|$ for all $u$. We also report the value obtained by a (full) fast CBC search for a PLR. The results are displayed in the first panel of Figure 1. We see that the FOM distribution has a smaller mean and much less variance for the Sobol’ points than for the other constructions. Even the median obtained for Sobol’ beats (slightly) the FOM obtained by a full CBC construction with PLRs. The quantiles for random PLRs and random nets are approximately the same.

The second panel of the figure shows the results of a similar sampling for PLRs with $\mathcal{I}_2$ as a FOM, also in 6 dimensions. Note that this FOM applies only to PLRs. Here, the FOM values are more dispersed and the fast CBC gives a significantly better value than the best FOM obtained by random sampling. Also the search for the point set parameters is much quicker with the fast CBC construction than with random sampling of size 1000.

Fig. 1: The 0.1, 0.5, and 0.9 quantiles of the FOM distribution as functions of $n$ for various constructions, in log-log scale.
6.2 Comparison with tabulated parameter selections

We now give small examples showing how searching for custom parameter values with LatNet Builder can make a difference in the RQMC variance compared with pre-tabulated parameter values available in software or over the Internet. We do this for Sobol’ nets, and our comparison is with the precomputed direction numbers obtained by Joe and Kuo [24], which are arguably the best proposed values so far. These parameters were obtained by optimizing a FOM based on the \(t\)-values over two-dimensional projections, using a CBC construction. With LatNet Builder, we can account for any selected projections in our FOM. For instance, if we think all the projections in two and three dimensions are important, we can select a FOM that accounts for all these projections. To illustrate this, we made a CBC construction of \(n = 2^{12}\) Sobol’ points in \(s = 15\) dimensions, using the sum or the maximum of \(t\)-values in the two- and three-dimensional projections. Figure 2 shows the distribution of \(t\)-values obtained with the sum, the max, and the points from [24]. Compared with the latter, we are able to reduce the worse \(t\)-value over 3-dim projections from 8 to 5 when using the max, and to reduce the average \(t\)-value when using the sum. However, when using the max, we get a few poor two-dim projections, because we compare the \(t\)-values on the same scale for two and three dimensions. We should probably multiply the \(t\)-value by a scaling factor that decreases with the dimension.

Fig. 2: Distributions of \(t\)-values for 2-dim and 3-dim projections, for three Sobol’ point sets: (1) Joe-Kuo taken from [24], (2) Max and (3) Sum are found by LatNet Builder as explained in the text. For each case, we report the number of projections having any given \(t\)-value, as well as the average \(t\)-value (dashed vertical lines).

In our next illustration, we compare the RQMC variances for Sobol’ points with direction numbers taken from [24] and direction numbers found by LatNet Builder using a custom FOM for our function. We want to integrate

\[
 f(u) = \prod_{j=1}^{5} (\psi(u_j) - \mu) + \prod_{j=6}^{10} (\psi(u_j) - \mu),
\]
where \( \psi(u) = ((u - 0.5)^2 + 0.05)^{-1} \) and \( \mu = \mathbb{E}[\psi(U)] \approx 10.3 \) when \( U \sim U(0,1) \). This function is the sum of two five-dimensional ANOVA terms for a more general function taken from [11]. A good FOM for this function should focus mainly on these two five-dim projections, namely \( u = \{1, 2, 3, 4, 5\} \) and \( u = \{6, 7, 8, 9, 10\} \), and not on the two-dim projections as in [24]. So we made a search with the \( \mathcal{P}_2 \) criterion with weights \( \gamma_u = 1 \) for these two projections and 0 elsewhere, to obtain new direction numbers for \( n = 2^{20} \) Sobol’ points in 10 dimensions. Then we estimated the variance of the sample RQMC average over these \( n \) points with the two choices of direction numbers (those of [24] and ours), using \( m = 200 \) independent replications of an RQMC scheme that used only a random digital shift. The empirical variance with our custom points was smaller by a factor of more than 18.

### 6.3 Variance for another toy function

Here we consider a family of test functions similar to those in [52]:

\[
  f_{s,c}(u) = \prod_{j=1}^{s} \left( 1 + c_j \cdot (u_j - 1/2) \right)
\]

for \( u \in (0,1)^s \), where \( c = (c_1, \ldots, c_s) \in (0,1)^s \). The ANOVA components are, for all \( u \subset \{1, \ldots, s\} \),

\[
  (f_{s,c})_u(u) = \prod_{j \in u} c_j \cdot (u_j - 1/2),
\]

For an experiment, we take arbitrarily \( s = 3 \) and \( c = (0.7, 0.2, 0.5) \). We use LatNet Builder to search for good PLRs with a fast CBC construction, with product weights \( \gamma_j = c_j \), with the FOMs \( \mathcal{P}_2 \), \( \mathcal{I}_2 \), and \( \mathcal{I}_3 \) (whose interlacing factors \( d \) are 1, 2, and 3, respectively). For each \( n = 2^k \), \( k = 5, \ldots, 18 \), we estimate the RQMC variance with \( m \) independent replications of the randomization scheme, with \( m = 1000 \) for LMS+shift, and \( m = 100 \) for NUS. For the interlaced points, the randomization is performed before the interlacing, as in [13]. Figure 3 shows the variance as a function of \( n \), in log-log scale. We see that the two randomization schemes give approximately the same variance. However, the time to generate and randomize the points is much larger for NUS than for LMS+shift: around 10 times longer for \( 2^{11} \) points and 50 times longer for \( 2^{18} \) points. As expected, the variance reduction and the convergence rate are larger when the interlacing factor increases, although the curves are more noisy.
Fig. 3: Variance as a function of $n$ in log-log scale, for PLRs with two randomization schemes and three interlacing factors $d$, found with LatNet Builder. We also report the average time to generate and randomize the points.

| Nb. of points | Constr. time |
|---------------|--------------|
| 128           | 0.02 s       |
| 512           | 0.04 s       |
| 2048          | 0.12 s       |
| 8192          | 0.50 s       |
| 32768         | 2.25 s       |
| 131072        | 8.32 s       |

7 Conclusion

LatNet Builder is both a tool for researchers to study the properties of highly uniform point sets and associated figures of merit, and for practitioners who want to find good parameters for a specific task. It is relatively easy to incorporate new FOMs into the software, especially if they are in the kernel form (9).

Many questions remain open regarding the roles of the construction, the search method, the randomization, and (perhaps more importantly) the choice of the weights. It is our hope that the software presented here will spur interest into these issues.

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