Abstract—In this paper, we consider the problem of synthesizing a controller in the presence of uncertainty such that the resulting closed-loop system satisfies certain hard constraints while optimizing certain (soft) performance objectives. We assume that the hard constraints encoding safety or mission-critical specifications are expressed using Signal Temporal Logic (STL), while performance is quantified using standard cost functions on system trajectories. To ensure satisfaction of the STL constraints, we algorithmically obtain control barrier functions (CBFs) from the STL specifications. We model controllers as neural networks (NNs) and provide an algorithm to train the NN parameters to simultaneously optimize the performance objectives while satisfying the CBF conditions (with a user-specified robustness margin). We evaluate the risk incurred by the trade-off between the robustness margin of the system and its performance using the formalism of risk measures. We demonstrate our approach on challenging nonlinear control examples such as quadcopter motion planning and a unicycle.

I. INTRODUCTION

Safety-critical cyber-physical systems typically have hard safety specifications that must be met by all system behaviors. Additionally, system designers often specify performance objectives to address efficiency concerns, and seek controllers to optimize these objectives. For example, consider an autonomous vehicle (AV) following another vehicle. Here, the AV must satisfy the safety specification of maintaining a minimum safe distance ($d_{safe}$) from the lead vehicle, while also minimizing the travel time for the AV. Clearly, the vehicle can be safe with a high robustness margin by driving slower than required (maintaining distance much greater than $d_{safe}$), but this leads to sub-optimal performance w.r.t. the travel time objective. Simultaneously designing for safety and performance requires the designers to never violate safety requirements in favor of performance; however, it may be possible to trade-off the safety margin to achieve better performance. This generates some risk from the safety perspective: how risky is it to use a controller with a lower safety margin? that may perform better? In this paper, we systematically study this problem.

We assume that safety is specified using a real-time temporal logic such as Signal Temporal Logic (STL) [1]. STL has emerged as a powerful specification language in the various cyber-physical system applications [2]–[5]. STL properties are built from atomic subformulate that are predicates over real-valued signals, combined with Boolean logic connectives such as and ($\land$), or ($\lor$), and temporal logic operators such as eventually (F), always (G), and until (U) indexed by time intervals. For example consider a design objective for a quadcopter: “The quadcopter must rendezvous in one of two designated regions $R_1$ or $R_2$ 5 to 7 minutes after takeoff before getting as close as possible to a given target destination within 20 mins, while avoiding no-fly zones.” Let $p(\cdot)$ denote the position of the quadcopter. The hard safety specifications in this objective can be expressed by the following STL formula:

$$\varphi_{qc} \equiv F_{[5,7]}(p \in R_1 \lor p \in R_2) \land G_{[0,20]}(p \notin R_{nofly})$$ (1)

The soft specification requires us to minimize $d(p, p_{target})$, where $d$ is a distance function and $p_{target}$ is a fixed target location. The quantitative semantics of STL allow us to measure the degree of satisfaction of an STL property using the notion of a robustness value [6]. We say that a system has safety robustness margin $\rho^s > 0$ if the minimum robustness value across all system behaviors exceeds $\rho^s$. In our framework, a performance objective can be any differentiable, real-valued function of the system behavior.

There has been considerable amount of research on the problem of synthesizing controllers that guarantee STL specifications. For example, using approaches from motion planning [7], [8], model predictive control [3], [9], [10], reactive synthesis [11], [12], reinforcement learning [13], [14], imitation learning [15], [16], and through the use of control barrier functions [17], [18]. Of these approaches, the most relevant to our paper is the one based on using control barrier functions (CBFs) [19]. Let $s_k$ denote the state of the system at time $k$. A CBF describes a set $C$ such that for all system states $s_k \in C$, there exists a control action that ensures that $s_{k+1} \in C$. Control synthesis from CBFs has seen a lot of success in to training safe controllers in a variety of applications [19]–[21]. Recent work has focused on CBFs that provide more general classes of invariants such as timed reachability [22], [23] and fragments of STL [24]. Prima facie, synthesis of controllers to satisfy STL specifications may look like a well-studied problem, however, several open problems remain:

1) Existing work uses hand-crafted CBFs over limited fragments of STL [17], for example disjunctive STL specifications (see Eq. (1)) may not be addressed.

2) Existing approaches do not explicitly consider the trade-off between safety and performance. A naïve encoding of the problem requires using several Lagrange multipliers,
which may not scale in general.

3) Many existing approaches focus on control of linear systems or simple nonlinear systems.

4) Existing work does not quantify risk awareness in trading off the safety margin versus system performance.

To address the above challenges, we propose a framework based on three main ideas: (1) control barrier functions that encode hard safety constraints, (2) neural network (NN) controllers to control arbitrary nonlinear systems, (3) a training algorithm to learn the NN parameters that trades off between safety and performance. Our training algorithm for NN controllers uses a stochastic gradient optimization method that uses adaptive moments. A crucial aspect of our algorithm is to explicitly guide the search for NN parameters using a robustness margin parameter: across iterations, the optimizer alternates between satisfying safety and performance based on the robustness of the NN controller vis-à-vis the desired robustness margin.

Finally, we evaluate the risk-awareness for each controller by picking different robustness margins as design parameters. For this analysis, we utilize the recently formulated risk-aware verification approach [25] that uses risk measures such as value-at-risk and conditional-value-at-risk. We demonstrate the efficacy of our method on several examples of nonlinear systems and disjunctive STL safety specifications.

II. BACKGROUND

In this section, we provide the mathematical notation and the overall problem definition. We use bold letters to indicate vectors and vector-valued functions, and calligraphic letters to denote sets. Let $s$ and $a$ respectively be the variables denoting state and control inputs taking values from compact sets $S \subseteq \mathbb{R}^n$ and $a \subseteq \mathbb{R}^m$, respectively. We use the words action and control input interchangeably. We consider discrete-time nonlinear feedback control systems of the following form:\footnote{Our technique can handle continuous-time nonlinear systems as well. This requires zero-order hold discretization of the dynamics in a sound way to account for system behavior between sample times.}

$$s_{k+1} = f(s_k, a_k).$$

Here, $s_k$ and $a_k$ denote the values of the state and action variables at time $k$. We assume that the controller can be expressed as a parameterized function $\pi_\theta$, where $\theta$ is a vector of parameters that takes values in $\Theta$. Later in the paper, we instantiate the specific parametric form using a neural network for the controller. Given a fixed vector of parameters $\theta$, the parametric control policy $\pi_\theta$ returns an action $a_k$ as a function of the current state $s_k \in S$ and time $k \in \mathbb{Z}_{\geq 0}$. Namely,

$$a_k = \pi_\theta(s_k, k)$$

We will be using the terms controller and control policy interchangeably. Under a fixed policy, Eq. (2) is an autonomous discrete-time dynamical system. For a given initial state $s_0 \in I \subseteq S$ and dynamics $f$, a system trajectory $\sigma_{s_0, f}$ is a function from $[0, K] \subset \mathbb{Z}_{\geq 0}$ to $S$, where $\sigma_{s_0, f}(0) = s_0$, and for all $k \in [0, K - 1]$, $\sigma_{s_0, f}(k + 1) = f(s_k, \pi_\theta(s_k, k))$. To address modeling inaccuracies, we also consider bounded uncertainty in the model. We denote $\mathcal{F}$ as the family of possible realizations of the model $f \in \mathcal{F}$. If the policy $\pi_\theta$ is obvious from the context, we drop the $\theta$ in the notation $\sigma_{s_0, f}$.

The main objective of this paper is to formulate algorithms to obtain the optimal policy $\pi_\theta$ that guarantees the satisfaction of certain task objectives and safety constraints while optimizing performance rewards. In the rest of the section, we formulate controller synthesis as an optimization problem that we seek to solve. In order to define this formally, we first introduce a performance reward, and then introduce task objectives/safety constraints.

Performance reward. In practical control applications, it is common to quantify the control performance using a state-based reward function [26], [27]. Formally,

$$\mathcal{J}^{\text{perf}}(s_0, f, \theta) = \sum_{k=0}^{K} \gamma^k r(\rho_{s_0, f}(k))$$

Task Objectives and Safety Constraints. We assume that task objectives or safety constraints of the system are specified in a temporal logic known as Signal Temporal Logic (STL) [1]. STL formulas are defined using the following syntax:

$$\varphi = h(s) \triangleright \triangleright 0 \mid \varphi_1 \land \varphi_2 \mid \varphi_1 \lor \varphi_2 \mid F \varphi \mid G \varphi \mid \varphi_1 U \varphi_2$$

Here, $\triangleright \triangleright \in \{\leq, <, >, \geq\}$, $h$ is a function from $S$ to $\mathbb{R}$, and $I$ is a closed interval $[a, b] \subseteq [0, K]$.

Semantics. The formal semantics of STL over discrete-time trajectories have been previously discussed in [6]. We denote the formula $\varphi$ being true at time $k$ in trajectory $\sigma_{s_0, f}$ by $\sigma_{s_0, f}, k \models \varphi$. We say that $\sigma_{s_0, f}, k \models h(s) \triangleright \triangleright 0$ iff $h(\sigma_{s_0, f}(k)) \triangleright \triangleright 0$. The semantics of the Boolean operations ($\land$, $\lor$) follow standard logical semantics of conjunctions and disjunctions respectively. For temporal operators, we say $\sigma_{s_0, f}, k \models F \varphi$ is true if there is a time $k \in I$ where $\varphi$ is true. Similarly, $\sigma_{s_0, f}, k \models G \varphi$ is true iff $\varphi$ is true for all $k \in I$. Finally, $\sigma_{s_0, f}, k \models \varphi_1 U \varphi_2$ if there is a time $k \in I$ where $\varphi_2$ is true and for all times $k' \in [0, k)$ $\varphi_1$ is true.

In addition to the Boolean satisfaction semantics, STL also permits quantitative satisfaction semantics. These are defined with a robustness function $\rho$ evaluated over a trajectory. We omit the formal definition; it can be found in [1], [6]. Intuitively, the robustness function defines robustness of predicates at a given time $k$ to be proportional to the signed distance of the state variable value at $k$ from the set of values satisfying the predicate. Conjunctions and disjunctions map to minima and maxima of the robustness of their subformulas respectively. Temporal operators can be viewed as conjunctions/disjunctions (or their combinations) over time. We denote $\rho_{\varphi}(s_0, f)$ as the robustness of the
trajectory starting in state \( s_0 \) for the dynamics \( f \). Note that if \( \rho_\varphi(s_0, f) > 0 \) then it implies that \( \sigma_{s_0,f} \models \varphi \).

**Risk Measures.** We now introduce two commonly used risk measures that are used to quantify the risk of design mistakes in stochastic systems. Recall that \( F \) represents a family of possible realizations of the system dynamics; we assume that we are provided a probability distribution \( D_F \) on \( \mathcal{F} \), and \( D_Z \), a distribution over the initial states of the system. A risk measure maps a scalar random variable to the extended set of reals \( \mathbb{R} \cup \{ -\infty, \infty \} \). Typically, the input to the risk measure indicates a cost that we wish to optimize, and the risk measure quantifies the level of risk associated with a possibly sub-optimal result from the optimization. There exist various risk measures. We are particularly interested in tail risk measures (Value-at-Risk (VaR), Conditional-Value-at-Risk (CVaR)) [28] For a confidence probability \( \varepsilon \) and real-valued random variable \( Z \),

\[
VaR_\varepsilon(Z) = \inf_{\xi \in \mathbb{R}} \{ \xi \mid \Pr(Z \leq \xi) \geq \varepsilon \}
\]

(6)

CVaR\(_\varepsilon\)(\(Z\)) := \( \mathbb{E}[Z \mid Z \geq VaR_\varepsilon(Z)] \)

(7)

In the above definitions, if we consider the cost \( Z = -\rho_\varphi(s_0, f) \), and a given confidence probability \( \varepsilon \) to compute the corresponding risk measure \( r_\varepsilon \).

Essentially, both risk measures provide a quantification on risk by computing probabilistic upper bounds on the negative of the robustness value, or lower bounds on the actual robustness value, as required in risk-aware verification [25], [28].

**Problem Definition.** (i) Learn an optimal policy \( \pi_\theta(s_k, k) \) such that it satisfies a given STL formula \( \varphi \) while maximizing the performance reward defined in Eq. (4).

\[
\theta^* = \underset{\theta \in \Theta}{\arg \max} \quad \frac{\mathbb{E}}{s_0, f \sim (D_F \times D_Z)} \left[ \int \pi_{\text{perf}}(s_0, f, \theta) \right] \\
\text{s.t.} \quad \forall s_0 \in \mathcal{I}, \forall f \in \mathcal{F} : \sigma_{s_0,f} \models \varphi
\]

(8)

(ii) Given a confidence threshold \( \varepsilon \), we will quantify the level of risk \( r_\varepsilon \) according to the risk measures (6), (7).

**III. CONTROL BARRIER FUNCTIONS FOR STL**

In [17], the authors introduce time-varying control barrier functions that are used to synthesize controllers that are guaranteed to satisfy a given STL specification. We first adapt this notion to discrete-time nonlinear systems.

**Definition 3 (Discrete-Time Time-Varying Valid Control Barrier Functions (DT-CBF))** Let \( b : S \times [0, K] \to \mathbb{R} \) be a function that maps a state and a time instant to a real value. Let \( B(k) = \{ s_k \mid b(s_k, k) \geq 0 \} \) be a time-varying set. The function \( b \) is a valid, discrete-time, time-varying CBF if the zero level sets of the CBF are an envelope for any system trajectory, i.e.,

\[
\forall s_0 \in \mathcal{I}, \forall f \in \mathcal{F} : \forall k \in [0, K] : \sigma_{s_0,f}(k) \in B(k)
\]

(9)

**Algorithm 1: Recursive formulation of CBFs based on an STL formula**

```plaintext
1: Function stl2cbf(\( \varphi, \sigma_{s_0,k} \))
2: \text{case } \varphi = h(s, k) \geq 0
3: \quad \text{return } \mu(h(\sigma_{s_0,k})))
4: \text{case } \varphi = \varphi_1 \land \varphi_2
5: \quad \text{return } \text{softmin} \left( \text{stl2cbf}(\varphi_1, \sigma_{s_0,k}), \text{stl2cbf}(\varphi_2, \sigma_{s_0,k}) \right) \eta
6: \text{case } \varphi = \varphi_1 \lor \varphi_2
7: \quad \text{return } \text{wavg}(\text{stl2cbf}(\varphi_1, \sigma_{s_0,k}), \text{stl2cbf}(\varphi_2, \sigma_{s_0,k}), \beta^2)
8: \text{case } \varphi = \mathcal{G}_{[a,b]}^\varphi
9: \quad \text{return } \text{softmin} \left( \text{stl2cbf}(\varphi, \sigma_{s_0,k}) \eta \right)
10: \text{case } \varphi = \mathcal{F}_{[a,b]}^\varphi
11: \quad \text{return } \text{wavg}(\text{stl2cbf}(\varphi, \sigma_{s_0,k}), \beta^2)
12: \text{case } \varphi = \varphi_1 \mathcal{U}_{[a,b]} \varphi_2
13: \quad e \leftarrow -\infty
14: \text{for } k \leftarrow a \text{ to } b \text{ do}
15: \quad \quad e \leftarrow \text{wavg}(e, \text{stl2cbf}(\varphi_2, \sigma_{s_0,k}), \beta^2)
16: \text{return } e
```

**DT-CBF for STL.** We formulate CBFs in a recursive fashion based on the formula structure. We describe the overall procedure in Algorithm 1. Before we describe the actual algorithm, we introduce some helper functions. The softmin function defined in Eq. (10) has been used in the past by several approaches [16], [29], [30] as a smooth approximation for computing the minimum of a number of real-valued quantities. In our softmin function, similar to [29], we introduce an additional parameter \( \eta > 1 \) that is used to control the level of conservatism in the approximation. Intuitively, larger values of \( \eta \) reduce the conservatism in the minimum value and hence allow for a better representation of the safety margin. Later, we discuss how the softmin function appears as a part of a cost function to be optimized; we include \( \eta \) as a part of this optimization process.

\[
\text{softmin}(v_1, \ldots, v_k; \eta) = -\frac{1}{\eta} \ln \left( \sum_{i=1}^{k} e^{-\eta v_i} \right)
\]

(10)

We also define the weighted average function \( \text{wavg} \). We are interested in two different form of \( \text{wavg} \) presented in Eq. (11) and Eq. (12); here \( \beta = (\beta_1, \ldots, \beta_k) \).

\[
\text{wavg}_1(v_1, \ldots, v_k; \beta) = \sum_{i=1}^{k} \left( \frac{\beta_i^2}{\sum_{i=1}^{k} \beta_i^2} \right) v_i,
\]

(11)

\[
\text{wavg}_2(v_1, \ldots, v_k; \beta) = \sum_{i=1}^{k} \left( \frac{\exp(\beta_i)}{\sum_{i=1}^{k} \exp(\beta_i)} \right) v_i,
\]

(12)

where the former is more accurate but the latter is more efficient for gradient descent. Finally, we articulate useful properties of softmin and \( \text{wavg} \) in Lemma 3.1.
Lemma 3.1: For all $v_1, \ldots, v_k \in \mathbb{R}$, and for $\eta \in \mathbb{R}_{>1}$, the following are true:

$$(\min_i v_i) \geq \text{softmin}(v_1, \ldots, v_k; \eta)$$

$$(\max_i v_i) \leq \text{wavg}(v_1, \ldots, v_k; \beta) \leq (\max_i v_i)$$

We can now describe Algorithm 1. The function $b_\phi$ computes the CBF w.r.t. either an atomic signal predicate or a conjunction of atomic predicates. For convenience, we represent $\sigma_{s0,f}$ with $\sigma_{s0}$ as short-hand. The CBF for an atomic predicate $\phi$ of the form $h(s_k) > 0$ is defined using a function $\mu$ that ensures that $\mu(h(s_k))$ is positive if $h(s_k)$ is positive, 0 if it is zero and negative otherwise. The CBF of the conjunction of two predicates is simply the softmin of the CBFs of the conjuncts. In the function $\text{st12cbf}(\varphi, \sigma_{s0,f}, k)$, we consider four cases. If $\varphi$ is a formula of the form $G_{[a,b)}\phi$, then we return the softmin of $b_\phi(s_\ell, f)$ for all $\ell \in [a,b]$. If $\varphi$ is of the form $F_{[a,b)}\phi$, then we return the weighted average of the CBFs at all time instants in $[a,b]$. For conjunctions of either kinds of temporal formulas, we again return the softmin and for disjunctions, we return the weighted average. Note that the function $\text{st12cbf}$ can be invoked with a concrete trajectory whereupon it returns a numeric value. It can also be invoked with a symbolic trajectory (where the symbols $s_k$ indicate the symbolic state at time $k$), whereupon it returns a numeric value. The CBF is the smooth candidate CBF that is the smooth robustness of the trajectory and is a guaranteed lower bound for trajectory robustness $\rho_{s0}(s_0, f)$. 

Lemma 3.2: For any formula $\varphi$ belonging to STL, for a given trajectory $\sigma_{s0,f} = s_{0}, s_{1}, \ldots, s_{n}$, if $\text{st12cbf}(\varphi, \sigma_{s0,f}, 0) > 0$, then $\sigma_{s0,f} \models \varphi$.

Proof: We can prove this recursively over the formula structure and from the identities in Lemma (3.1). It is necessary to mention if $\text{st12cbf}(\varphi) < 0$ it does not imply the STL specifications are violated.

Example 1: Consider the STL specification in Eq. (15). 

$$F_{[1,10]}(s \in E_1) \vee F_{[1,10]}(s \in E_2) \wedge G_{[1,20]}(s \in \neg E_3)$$

Let $c_2 = (2,8), c_1 = (5,5), c_3 = (8,2)$, and $r = \sqrt{1.5}$, $\mu(s_k) \in [1,2]$. Then, in Eq. (15), for $i \in [1,3]$: $E_i = (s - c_i)^T(s - c_i) \leq r$. We define the CBF $\mu(s_k) \in [1,2]$ to be $1 - e^{-(s_k - c_i)^T(s_k - c_i) - r}$. For $j = 2,3$, we define $\mu(s_k) \in [1,2]$ to be $r - (s_k - c_j)^T(s_k - c_j)$. Then, the CBF w.r.t. the formula can be computed using Algorithm 1. 

IV. LEARNING-BASED CONTROL SYNTHESIS

We remark that the trajectory $\sigma_{s0,f}$ is essentially a repeated composition of f and the neural controller. Thus, we can compute the gradient of the performance costs and STL objectives (as expressed by the CBF) with respect to the controller parameters $\theta$ using standard backpropagation methods for neural networks.

A. Training Neural Networks to satisfy specifications

We explain the procedure for training a neural controller w.r.t. performance and safety specifications in algorithm 2. The training algorithm aims to approximate the solution of Eq. (8). Thus, the first step is to reformulate it free from constraints. Algorithm 1 provides the smooth trajectory robustness for a given STL formula $\varphi$. This robustness is a function of the common variable $\eta$ that is used by all softmax functions, and the tuple of $\beta$ variables for each subformula $\psi$ of a disjunctive formula (denoted $\beta^\psi$). In addition to the neural network parameters $\theta$, we also treat the variables $\eta$ and the sets $\{\beta^\psi\}$ as decision variables in the training process. The $\beta$ variables are already unconstrained; however, $\eta$ is constrained: $\eta > 1$. To remove this constraint, we introduce $\eta = \lambda^2 + 1$, and use $\lambda$ as a decision variable. We denote the tuple of conjunctive variable ($\lambda$) and all disjunctive variables ($\beta^\psi$) with $\nu$. We also denote this robustness with $J_{STL}^\nu$ that is a function of tuple $(s_0, f, \theta, \nu)$.

$$J_{STL}(s_0, f, \theta, \nu) = \text{st12cbf}(\varphi, \sigma_{s0,f}, 0), \quad \nu = (\{\beta^\psi\}, \lambda).$$

We also sample a batch $\tilde{I}$ of initial states uniformly from $\tilde{I}$ for training purposes.

The training algorithm is primarily inspired by the Lagrange multiplier technique that transforms a constrained optimization to non-constrained,

$$J = \max_{\theta, \nu} \sum_{s_0 \in \tilde{I}, f \in F} (J_{perf}(s_0, f, \theta) + \omega_{s0,f}J_{STL}(s_0, f, \theta, \nu)).$$

First order optimality conditions guarantee that as long as the Lagrange multipliers are positive, the cost as defined using $J_{STL}$ is positive. This in turn guarantees the satisfaction of STL specifications along the trajectory. Since $J_{STL}(s_0, f, \theta, \nu)$ is highly non-convex, optimization (17) is quite intractable and the solution may not satisfy the KKT optimality condition [31, [32]. However, the main role of the Lagrange multipliers $\omega_{s0,f}$ is to perform a trade-off between $J_{STL}$ and $J_{perf}$ and one of the contributions of this work is to propose a training process that focuses on applying this trade off. The training algorithm utilizes the gradients $\nabla_{\theta, \nu}J_{perf}(s_0, f, \theta)$ and $\nabla_{\theta, \nu}J_{STL}(s_0, f, \theta, \nu)$ (obtainable from Automatic differentiation package [33]). In case, the cost specified with $J_{STL}$ is less than a user-specified threshold, then the algorithm increases this with an adaptive selection between $\nabla_{\theta, \nu}J_{perf}(s_0, f, \theta)$ and $\nabla_{\theta, \nu}J_{STL}(s_0, f, \theta, \nu)$. Otherwise, it increases the performance with $\nabla_{\theta, \nu}J_{perf}(s_0, f, \theta)$. We call the mentioned user specified threshold as robustness margin (denoted by $p$).

We now describe Algorithm 2. We use the variable $i$ to denote the iteration number during training. We use the notation $(\theta, \nu)^i$ to denote the value of $\theta$ and $\nu$ at the beginning of iteration $i$. We initialize $(\theta, \nu)^0$ randomly. At the beginning of each training iteration, in line 3 we sample $f$ from $F$, then in lines (4-6), for all states in $\tilde{I}$, we calculate the gradients $\nabla_{\theta}J_{perf}(s_0, f, \theta, \nu), \nabla_{\nu}J_{perf}(s_0, f, \theta)$ (stored in $d_1(s_0)$), and $\nabla_{\theta}J_{STL}(s_0, f, \theta, \nu) \nabla_{\nu}J_{STL}(s_0, f, \theta, \nu)$ (stored in $d_2(s_0)$). Of these, note that $\nabla_{\nu}J_{perf}(s_0, f, \theta)$ is 0. We then compute the state $b_1$ (resp. $b_2$) for which the 2-norm of $d_1(s_0)$ (resp. $d_2(s_0)$) is the highest. The highest values of the states $b_1$ and $b_2$ are respectively stored in $d_1$ and $d_2$ (Line 8).

The next step is to compute potential updates to the parameter values $\theta$ and the STL parameters $\nu$ (Lines 9-11). Roughly, the values $(\theta_1, \nu_1)$ represent the update to $(\theta, \nu)^i$ using only the inclusion of gradient for performance cost in
Algorithm 2: Sampling based algorithm for training the parameterized policy.

1. \( i \leftarrow 0 \), Initialize \((\theta, v)\)\(^0\), Sample \( I \) to obtain \( \tilde{I} \)
2. while true do
3. \( \text{Sample } f \in F \)
4. foreach \( s_0 \in \tilde{I} \) do
5. \[
\delta_1(s_0) \leftarrow [\nabla_\theta J_{\text{perf}}(s_0, f, \theta), 0] \\
\delta_2(s_0) \leftarrow [\nabla_\theta J_{\text{STL}}(s_0, f, \theta, v), \nabla_v J_{\text{STL}}(s_0, f, \theta, v)]
\]
6. // get states with the best grad. values
7. \( b_1, b_2 \leftarrow \arg\max ||\delta_1(s_0)||_2, \arg\max ||\delta_2(s_0)||_2 \)
8. \( s_0 \in \tilde{I} \)
9. \( d_1, d_2 \leftarrow \delta_1(b_1), \delta_2(b_2) \)
10. // candidate parameter updates
11. \( (\theta_1, v_1) \leftarrow (\theta, v)^i + \text{Adam}(d_1) \)
12. \( (\theta_{\text{STL}}, v_{\text{STL}}) \leftarrow (\theta, v)^i + \text{Adam}(d_2) \)
13. \( (\theta_{\text{slow}}, v_{\text{slow}}) \leftarrow (\theta, v)^i + \text{Adam}(d_1)/\tau \)
14. Sample \( s_i \) from \( I \)
15. /* Pick update giving best tradeoff between perf. and safety */
16. if \( J_{\text{STL}}(s_0, f, (\theta, v)^i) \leq \rho \) then
17. if \( J_{\text{STL}}(s_0, f, \theta_1, v_1) \geq J_{\text{STL}}(s_0, f, (\theta, v)^i) \) then
18. \( (\theta, v)^{i+1} \leftarrow (\theta_1, v_1) \)
19. else
20. \( (\theta, v)^{i+1} \leftarrow (\theta_{\text{STL}}, v_{\text{STL}}) \)
21. else
22. \( (\theta, v)^{i+1} \leftarrow (\theta_{\text{slow}}, v_{\text{slow}}) \)
23. if \( ||(\theta, v)^{i+1} - (\theta, v)^i||_2 \leq \varepsilon \) then
24. return \((\theta, v)^i\)
25. terminate

Adam optimizer. The values \((\theta_{\text{STL}}, v_{\text{STL}})\) represent the update only using the gradient of smooth trajectory robustness in Adam optimizer. Finally, \((\theta_{\text{slow}}, v_{\text{slow}})\) represents a slower update with gradient of performance for some \( \tau > 1 \).

Next, we sample a state \( s_i \) uniformly at random and use it for cost computation. If \( J_{\text{STL}}(s_i, f, (\theta, v)^i) < \rho \), i.e., our user-provided robustness margin (Line 13), then we need to take steps to increase the smooth trajectory robustness. We consider two cases: (1) If using the update based on the gradient of the performance cost improves the smooth trajectory robustness, we choose this update as it allows us to improve both performance and robustness, i.e., satisfaction robustness (Line 15). (2) Otherwise, we use the update based on the gradient of the smooth trajectory robustness \( J_{\text{STL}} \) (Line 17).

If \( J_{\text{STL}}(s_0, f, (\theta, v)^i) \geq \rho \), then we are robustly satisfying our STL constraints. In further quest to improve the performance cost, we need to take care that we do not reduce the robustness margin w.r.t. STL constraints. Hence, we use a slower learning rate that takes smaller steps in trying to improve the performance (Line 19).

Remark 1: Considering that this algorithm only focuses on increasing \( J_{\text{STL}} \) up to \( \rho \geq 0 \), once the STL specification is satisfied then it focuses on optimizing performance. In a sense, this switching strategy plays a role similar to that of Lagrange multipliers: performance cost is optimized only if the robustness is above the user-provided threshold.

B. Risk estimation

The minimum number of samples to guarantee the confidence on the verification results is proposed in [34]. We generate \( N = 10^6 \) samples \((s_0, f)\) uniformly from \((I \times F)\) and simulate the corresponding trajectories \( \sigma_{s_0,f} \). We compute the robustness \( \rho_s(s_0, f) \) for every single trajectory and calculate \( \text{VaR} \) through obtaining the \( \varepsilon * 100 \) percentile of the negation of the robustness values [35] and calculate \( CVaR \) according to definition 7.

V. EXPERIMENTAL EVALUATION

A. Unicycle Dynamics

We demonstrate the efficacy of our technique on a nonlinear unicycle model. We define the uncertainty for the initial condition as:

\[ I = \{ s_0 | (x_0, y_0, \alpha_0) \in [0.6, 1.4] \times [0.6, 1.4] \times \left[ \frac{2\pi}{3}, \frac{4\pi}{3} \right] \} \]

The unicycle dynamics with uncertainties are defined as follows,

\[
\begin{align*}
x_{k+1} &= \left(1 + \delta\right)x_k + v_k/\omega_k \left(\sin(\alpha_k + \omega_k) - \sin(\alpha_k)\right) \\
y_{k+1} &= \left(1 + \delta\right)y_k + v_k/\omega_k \left(\cos(\alpha_k) - \cos(\alpha_k + \omega_k)\right) \\
\alpha_{k+1} &= \left(1 + \delta\right)\alpha_k + \omega_k
\end{align*}
\]

where \( \delta \in [-0.01, 0.01] \) and the control inputs, \( v_k, \omega_k \) are bounded: \( v_k \in [0, 1] \), \( \omega_k \in [-0.5, 0.5] \). To restrict the controller in proposed bounds we fix the last hidden layer of neural controller [sigmoid, tanh] and include it to model. Thus, we reformulate the dynamics by replacing the controllers with:

\[
\begin{align*}
v_k &\leftarrow \text{sigmoid}(0.5\sigma_1(k)), \quad a_1(k) \in \mathbb{R} \\
\omega_k &\leftarrow 0.5 \tanh(0.5\sigma_2(k)), \quad a_2(k) \in \mathbb{R}
\end{align*}
\]

B. quadcopter Dynamics

In another attempt we consider controlling a quadcopter with uncertain dynamics. We define the uncertainty for the initial condition as a spherical set, \( I = B_r(c) \) with center, \( c = [0.025, 0.025, 0, 0, 0, 0]^T \) and radius \( r = 0.0125 \). The quadcopter also follows the following uncertain dynamics,

\[
\begin{align*}
x(k+1) &= \left(1 + \delta\right)x(k) + 0.05v_x(k) \\
y(k+1) &= \left(1 + \delta\right)y(k) + 0.05v_y(k) \\
z(k+1) &= \left(1 + \delta\right)z(k) + 0.05v_z(k) \\
v_x(k+1) &= (1 + \delta)v_x(k) + 0.4905\tan(u_1(k)) \\
v_y(k+1) &= (1 + \delta)v_y(k) - 0.4905\tan(u_2(k)) \\
v_z(k+1) &= (1 + \delta)v_z(k) + 0.05(g - u_3(k))
\end{align*}
\]

discretized with ZOH for timestep \( T = 0.05 \) sec. Here \( \delta \in [-0.01, 0.01] \) and the control inputs, \( u_1(k) \in \)
Fig. 1: Figures (a) and (b) present the evolution of performance cost (blue) vs trajectory robustness (green) over the training process of unicycle dynamics for $\rho = 0.5, 0.3$ respectively. The horizontal axis is presented in log form.

Fig. 2: Presents the evolution of performance cost (blue) vs trajectory robustness (green) over the training process in quadcopter example for $\rho = 0.1$. The horizontal axis is in log form.

| Example   | $\rho$ | $\tau$ | Controller Dimension | Activation Function | Iterations | Runtime (secs) | Expected value | Expected value |
|-----------|--------|--------|----------------------|--------------------|------------|----------------|----------------|----------------|
| Unicycle  | 0.3    | 1e2    | [4,5,2,2]            | tanh               | 40000      | 1048           | 35.3430        | 0.6108 / 0.6109|
| Unicycle  | 0.5    | 1e2    | [4,5,2,2]            | tanh               | 40000      | 1067           | 33.3528        | 0.6108 / 0.6109|
| quadcopter| 0.1    | 5e4    | [7,10,3,3]           | tanh               | 10000      | 155            | 24.3024        | 0.6729 / 0.7516|

TABLE I: Training and Validation Results

| Example   | CBF for atomic propositions, $h_a(s_0, k)$ | Reward | discount |
|-----------|--------------------------------------------|--------|----------|
| Unicycle  | $\phi_1: \sigma_a(k) \in \mathcal{E}_1, k \in [1, 10]$ | $10\exp\left(-\frac{(s_3 - 2)^2 + (y_3 - 2)^2}{0.01^2}\right)$ | 0.9 |
| quadcopter| $\phi_1: \sigma_a(k) \in \mathcal{E}_1, k \in [1, 10]$ | $10\exp\left(-\frac{(s_3 - 2)^2 + (y_3 - 2)^2}{0.01^2}\right)$ | 0.9 |

TABLE II: Shows the CBFs and reward functions we utilize in training process.

Fig. 3: (a) Represents sample trajectories with the random initial value for $\theta$, (b,c) respectively show sample trajectories for trained $\theta$ with robustness margin $\rho = 0.3$ and 0.5. This figure clearly shows the trajectories shift towards the center of $\mathcal{E}_1$ when the robustness margin $\rho$ increases. For this simulation we sample 500 different $(s_0, f)$ uniformly at random from $\mathcal{I} \times \mathcal{F}$ and simulate the trajectories. The green plots satisfy the STL specifications while its darkness shows the level of performance. There exists 2 red trajectories in (b) that are marginally violating the STL specs.

$$[-0.1, 0.1], \ u_2(k) \in [-0.1, 0.1], \ u_3(k) \in [7.81, 11.81].$$

The parameter $g = 9.81$ is the gravity. To impose bounds on the controller, like the Unicycle example, we fix the last hidden layer of the neural controller, \[
\begin{align*}
\text{tanh}, \ \text{tanh}, \ \text{tanh}
\end{align*}
\] and include it in the model,

\[
\begin{align*}
u_1(k) & \leftarrow 0.1 \text{tanh}(0.1a_1(k)), \quad a_1(k) \in \mathbb{R} \\
u_2(k) & \leftarrow 0.1 \text{tanh}(0.1a_2(k)), \quad a_2(k) \in \mathbb{R} \\
g - u_3(k) & \leftarrow 2 \text{tanh}(0.1a_3(k)), \quad a_3(k) \in \mathbb{R}
\end{align*}
\]
C. Results

The STL specifications for both examples are adopted from [15] and are introduced in Eq. (15) (Example 1). Regions $E_1, E_2$ and $E_3$ for unicycle and quadcopter examples are introduced in Fig 3 and Fig. 4 respectively. The unicycle and quadcopter approaches to the target $O = [8, 8]^\top$, $O = [0.1, 0.1, -0.0375]$ respectively. They are planned to approach $O$ with the highest possible level of performance (fast and close) within $K = 20$ time steps. The reward function and CBFs are defined in table II for both examples. We train a controller to satisfy the performance and STL task for unicycle and quadcopter dynamics. Table I shows the training result for $\rho = 0.3, 0.5$ in unicycle and $\rho = 0.1$ in quadcopter example. This table shows the trade-off between performance and STL robustness for the unicycle example.

We utilized (11) and (12) in training the disjunctive parameters $\beta$ for unicycle and quadcopter respectively. Fig. 5 shows the evolution of disjunctive parameters over the training process. Fig. 1 and 2 present the trade-off between the performance cost $J^{\text{perf}}$ and trajectory robustness $J^{\text{STL}}$ over the training process for both examples. Fig. 3 and Fig. 4 present the simulation of trajectories for unicycle and quadcopter examples, respectively.

Table III presents the results on probabilistic verification or risk-analysis for the controllers. For the unicycle dynamics, we can see that increasing the robustness margin parameter $\rho^*$ leads to an increase in the (probabilistic) lower bound on the robustness. Increasing the confidence level reduces the probabilistic lower bound. In fact, at 99.9% confidence, there is a risk of seeing system behaviors that violate the specifications by a margin of 0.133. Similar risks can be seen at the 99% and 99.9% confidence in the CVaR values. Intuitively, table III matches our expectation that controllers designed with higher robustness margin should have lower risk of violating specifications (at the cost of performance).

VI. Conclusion and Future Work

In this work we propose the weighted average, a useful tool to include disjunctive STL formula in the existent soft constrained policy optimization techniques [17]. We also utilize time dependent feedback policies that facilitates control in presence of STL specifications. This enables us to control the model with smaller neural networks. Non-convex optimizations may be intractable for Lagrange multiplier techniques. We address this problem with proposition of a training algorithm that simulates the trade off between objective and its constraints. We finally utilize this training algorithm for non-convex policy optimization with respect to STL specifications.

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TABLE III: Risk measures with one million data points.

| Confidence | Unicycle Dynamics | quadcopter Dynamics |
|------------|-------------------|--------------------|
| $\rho$ = 0.3 | $-\text{VaR}_0$, $-\text{CVaR}_0$ | $-\text{VaR}_0$, $-\text{CVaR}_0$ |
| $\rho$ = 0.5 | $0.246$, $0.132$ | $0.540$, $0.417$ |
| $\rho$ = 0.1 | $0.527$, $0.455$ | $0.527$, $0.455$ |
| $\epsilon$ | $0.059$, $0.027$ | $0.336$, $0.239$ |
| $0.95$ | $0.406$, $0.360$ | $0.406$, $0.360$ |
| $0.99$ | $0.305$, $0.284$ | $0.305$, $0.284$ |
| $0.999$ | $0.305$, $0.284$ | $0.305$, $0.284$ |

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Fig. 4: (a) Represents 500 trajectories generated with trained controller parameters $\theta$ for $\rho = 0.1$. For this simulation, we sample 500 different $(s_0, f)$ uniformly at random from $(I \times F)$ and simulate the trajectories. The darkness of trajectories is corresponding to their level of performance. There is no trajectory violating the STL specification. (b, c, d) shows the projection of trajectories on X-Y, Y-Z and X-Z planes respectively. (e) Represents simulated sampled trajectories for the initial value of $\theta$ that we utilized in training process.

Fig. 5: (a, b) shows the evolution of disjunctive parameters over the training process in unicycle ($\rho = 0.3$) and quadcopter ($\rho = 0.1$) examples, respectively. The log-scale horizontal axis indicates number of training iterations. There are three disjunctive formulas in (15): $F_{[1,10]}(s \in E_1)$, (that needs parameters $\beta_1, \ldots, \beta_{10}$) in its CBF, $F_{[1,10]}(s \in E_2)$ (using parameters $\beta_21$ and $\beta_22$ and the disjunction between these formulas that uses parameters $\beta_21$ and $\beta_22$. (a) Parameter $\beta_22$ converging to zero indicates that the system chooses to satisfy the first subformula thus the variables $\beta_1, \ldots, \beta_{10}$ are not relevant and not plotted. The $\beta_2$ parameter has the largest value, indicating the majority of the trajectories are in region $E_1$ at time $k = 7$. (b) Here, $\exp(\beta_21)$ converging to zero implies that $\beta_1, \ldots, \beta_{10}$ are not relevant. As the parameters $\beta_{19}, \beta_{20}$ are nonzero, the majority of trajectories are in $E_2$ at $k = 9$ and the others at $k = 10$.

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