Contextuality is one way of capturing the nonclassicality of quantum theory. The contextual nature of a theory is often witnessed via the violation of noncontextuality inequalities—certain linear inequalities involving probabilities of measurement events. Using the exclusivity graph approach (one of the two main graph theoretic approaches for studying contextuality), it was shown [Cabello et al. Phys. Rev. A 88, 032104 (2013); Chudnovsky et al. Ann. Math. 164, 51 (2006)] that a necessary and sufficient condition for witnessing contextuality is the presence of an odd number of events (greater than three) which are either cyclically or anticyclically exclusive. Thus, the noncontextuality inequalities the underlying exclusivity structure of which is as stated, either cyclic or anticyclic, are fundamental to quantum theory. We show that there is a unique noncontextuality inequality for each nontrivial cycle and anticycle. In addition to the foundational interest, we expect this to aid the understanding of contextuality as a resource to quantum computing and its applications to local self-testing.

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I. INTRODUCTION

A. Motivation

In an attempt to conceptually understand the departure of the predictions of quantum mechanics (QM) from that of classical physics, the notion of contextuality was introduced. It is one of the most general ways of capturing this divergence [1,2]; the celebrated Bell nonlocality can be viewed as a special case of contextuality where the context is provided via spacelike separation of the parties involved [3,4]. More generally, a context is defined by a set of compatible observables, viz., jointly measurable observables. Investigations into these fundamental questions have also reaped practical benefits. Bell nonlocality has found many applications in quantum key distribution [5], randomness certification [6], self-testing [7–9], and distributed computing [10], to name a few [11]. Recently, contextuality has also been applied more directly to quantum key distribution [12,13] and variants of randomness certification [14] and self-testing [15]. Further, it has been uncovered to be the resource powering the measurement based model and a class of fault tolerant models of quantum computation [16,17], among others [1,16–23].

B. Bell nonlocality or contextuality

The idea at the heart of this discussion can be traced back to Einstein, who expressed his discomfort with the probabilistic

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nature of quantum mechanics by providing a striking argument against it [24] using a notion of realism (element of physical reality) for two spatially separated experiments. He believed that there must exist local hidden variables which, once supplied, make QM deterministic. Such completions are referred to as local hidden variable models. Bell constructed a linear inequality which is violated by QM and yet it can never be violated by any such completion [3], falsifying Einstein’s belief [25]. It may be said that the Bell inequality witnesses the nonlocality of (any such completion) QM.

General discussions on this topic are facilitated by correspondingly considering general probabilistic assignments to the various observable events. The set of probabilistic assignments, which admit a local hidden variable description, form a convex polytope (a bounded set the boundaries of which are defined by hyperplanes). The facet-defining Bell inequalities are the characterizing hyperplanes of the aforesaid polytope. Once formalized, this becomes a general framework for studying Bell inequalities (which can and has been refined to facilitate computations). This can, however, be further generalized if an underlying principle which is correspondingly more general than that of local realism is used. In the Bell scenario, there was a clear role of spatial separation and therefore there were at least two parties involved. It turns out that one can study nonclassicality even for a single indivisible quantum system. To this end, one uses noncontextual completions of probabilistic assignments where the phrase noncontextual emphasizes that there is a precise value assigned to each observable by the completion. This is because it is possible to define completions where the value assigned depends on the context (i.e., the set of compatible observables it is measured with), and such completions can explain the predictions of quantum mechanics. Consequently quantum
mechanics is sometimes called contextual [26]. This schism between noncontextual completions and quantum mechanics is used as the underlying principle to construct frameworks to study noncontextuality (NC) inequalities. The Klyachko-Can-Binicioğlu-Shumovsky (KCBS) inequality may be considered to be the simplest NC inequality [35]. We clarify why in the exclusivity graph approach, an easy characterization of anticycles allows us to make a much more general statement about all scenarios (due to the strong perfect graph theorem, as was noted).

The study of the simplest Bell scenario using the exclusivity graph approach was carried out in [36] and the fundamental NC inequality was shown to be a Bell inequality involving only five events, termed a pentagonal Bell inequality (while the CHSH-Bell inequality involves eight events). However, the one involving seven events was missing.

### II. PRELIMINARIES

We summarize the exclusivity graph approach here, following the work of Amaral and Cunha [32], deferring a more complete discussion to the Appendix B. An outcome, \( a \), and its associated measurement, \( M \), are together called a measurement event (or events for brevity) and denoted by \((a|M)\). Two events are equivalent if their probability of occurrence is the same for all preparations. Let \( p_j(k) \) be the probability of getting an outcome \( k \) given that a measurement \( j \) was performed. Two events, \( e_i \) and \( e_j \), are exclusive if there exists a measurement \( M \) such that \( e_i \) and \( e_j \) correspond to different outcomes of \( M \) (see Definition 24 in Appendix B). With a family of events \( \{e_1, e_2 \ldots e_n\} \), we associate the exclusivity graph \( G := (V,E) \) where \( V \) is the set of vertices and \( E \) is that of edges, the vertices of which are the events, and there is an edge between the vertices if and only if the events are exclusive (see Definition 25 in Appendix B). The probabilities assigned to these events are formally given by a behaviour which for \( G \) is defined to be a map \( \map{p:V}{[0,1]} \) that assigns to each vertex \( i \) a probability \( p(i) \) such that \( p(i) + p(j) \leq 1 \) for all vertices that share an edge. The map \( p \) can also be seen as a vector in \( \mathbb{R}^{|V|} \) (see Definition 27 in Appendix B). Behaviours which admit a noncontextual completion, i.e., there exists a noncontextual hidden variable assignment such that if the hidden variable is traced out we recover the given behavior, are defined to be noncontextual behaviors (see Definition 28 in Appendix B). The set of such behaviors is denoted by \( \mathcal{B}_{NC}(G) \). We can similarly define the set of quantum behaviors, \( \mathcal{B}_Q \), to be that which can be obtained by at least one quantum state and corresponding observables (see Definition 29 in Appendix B). The set of E-principle behaviors, \( \mathcal{B}_E(G) \), is one where the behaviors respect the exclusivity principle (also referred to as the E principle), i.e., exclusive events must have their probability sum to at most one (see Definition 31 in Appendix B). The central claim of this formalism is that \( \mathcal{B}_{NC}(G) \subseteq \mathcal{B}_Q(G) \subseteq \mathcal{B}_E(G) \) (see Definition 31 in Appendix B). This is a corollary of a powerful identification of each of the sets with geometrical objects studied by Lovász which we describe later. We can now define more precisely a facet-defining NC inequality as being a nontrivial facet of \( \mathcal{B}_{NC}(G) \) where the direction of the inequality is chosen to satisfy containment in \( \mathcal{B}_{NC}(G) \) (see Definition 32 in Appendix B). An \( n \)-cycle graph is an \( n \) vertex graph where every \( i \)th vertex is connected to the \((i + 1)\)th vertex (the addition is modulo \( n \)). We define the fundamental cyclic noncontextuality (FCNC) inequality, corresponding to the \( n \)-cycle graph for \( n \) odd, to be

\[
K_n := \sum_{i=1}^{n} p_i \leq \frac{n-1}{2}.
\]

D. Relation to prior work

In the compatibility hypergraph approach, scenarios captured by odd \( n \)-cycle graphs are characterized by \( 2^n - 1 \) nontrivial NC inequalities, which include the generalized KCBS inequality [35]. We clarify why in the exclusivity graph approach, for odd \( n \)-cycle graphs, we obtain an exponential simplification—a unique NC inequality (see Sec. VII). The relevance of anticycles is not \textit{a priori} clear in the compatibility hypergraph approach and therefore, to the best of our knowledge, they have not been studied. However, in the exclusivity graph approach, an easy characterization of anticycles allows us to make a much more general statement about all scenarios (due to the strong perfect graph theorem, as was noted).
We analogously define the fundamental anticyclic noncontextuality inequality for the complement of the odd $n$-cycle graph as

$$K_0 := \sum_{i=1}^{n} p_i \leq 2.$$  

(2)

We recover the original KCBS inequality [37] in the special case of $n = 5$, both for the cyclic as well as the anticyclic case.

Relevant graph theoretic notions

We give a primer on the relevant graph theoretic notions which facilitates the statement and proof of our results.

Definition 1. The graph $G = (V, E)$ is defined by the set of vertices and the set of edges.

Definition 2. The orthonormal representation with respect to a graph $G$ is defined as follows. For all $i \rightarrow |v_i\rangle$ in $\mathbb{R}^d$, $(|v_i\rangle|v_j\rangle) = 0$ whenever $(i, j) \notin E$.

Definition 3. For a vector $|v_i\rangle$ in an orthonormal representation, the cost is defined as

$$c_i = ||(\psi|v_i\rangle||^2$$

where $|\psi\rangle \doteq (1, 0, \ldots, 0)$ is a vector in $\mathbb{R}^d$.

Definition 4. The theta body corresponding to a graph $G$ is defined to be

$$TH(G) = \{ p \in \mathbb{R}^{|V|} | p_{(i)} = c_i \}$$

where $c_i$ is the cost (see Definition 3) corresponding to $\tilde{G} := (V, \tilde{E})$ where $\tilde{E}$ is the complement [38] of $E$.

Definition 5. A stable set or independent set is a subset of vertices $K \subseteq V$ such that for all $i, j \in K$ there is no edge between $i$ and $j$. viz., $(i, j) \notin E$.

Definition 6. The independence number of a graph $G$ is defined to be the cardinality of the largest independent set of $G$.

Definition 7. A clique is a subset of vertices $K \subseteq V$ such that for all $i, j \in K$ there is an edge between $i$ and $j$, viz., $(i, j) \notin E$.

Definition 8. The incidence vector of a set is defined to be a vector $\vec{p}$ (of size $|V|$) for $K \subseteq V$ such that

$$p_{(i)} = \begin{cases} 1 & \text{if } i \in K \\ 0 & \text{else.} \end{cases}$$

Example 1. Consider the five-cycle graph $V = \{1, 2, 3, 4, 5\}$. $E = \{(1, 2), (2, 3), (3, 4), (4, 5), (5, 1)\}$. $K = \{1, 3\}$ is an example of a stable set. $K' = \{1, 2\}$ is an example of a clique. The incident vector corresponding to $K$ is $p = (1, 0, 1, 0, 0)^T$.

Definition 9. STAB($G$) (not to be confused with the stable set) is defined as the convex hull of the vectors $\vec{p}_{(k)}$ for all stable sets $k$ where $\vec{p}_{(k)}$ is the incidence vector of the set $k$. (Note: if $k$ were an index, $\vec{p}_{(k)}$ would refer to the kth component of the vector $\vec{p}$; here $k \in \mathbb{Z}$ is a set).

Definition 10. QSTAB($G$) is the set of vectors $x \in \mathbb{R}^{|V|}$ such that $x_i \geq 0$, $\sum_{i \in Q} x_i \leq 1$ for every clique $Q$.

Lemma 1 [39]. STAB($G$) is the convex hull of the integer solutions to the equations $x_i \geq 0$, $(x_i + x_j) \leq 1$ for every $(i, j) \in E$, where $G = (V, E)$.

Remark 1. Every set of indices which is an edge is also a clique (the other way is not necessary, obviously). This means that the inequalities listed in Definition 10 (viz., $\sum_{i \in Q} x_i \leq 1$ for every clique $Q$) contain the inequalities listed in Lemma 1 (viz., $x_i + x_j \leq 1$ for every $(i, j) \in E$).

Lemma 2 [40]. STAB($G$) $\subseteq$ TH($G$) $\subseteq$ QSTAB($G$).

III. UNIQUENESS OF FUNDAMENTAL CYCLIC NONCONTEXTUALITY INEQUALITIES

We begin by stating our first important result.

Theorem 1. Consider an odd $n$-cycle exclusivity graph. The associated FCNC inequality is a unique facet-defining NC inequality.

We will need the following powerful result connecting the behaviors to geometrically well-studied objects.

Lemma 3 [4]. Let $e_1, e_2, \ldots, e_n$ be the (exclusive) events with an associated exclusivity graph $G = (V, E)$. Then,

$$B_{NC}(G) = \text{STAB}(G),$$

$$B_{Q}(G) = \text{TH}(G),$$

$$B_{E}(G) = \text{QSTAB}(G).$$

Let QSTAB inequalities for a graph $G$ be the set of inequalities given by $\sum_{i \in Q} x_i \leq 1$ for every clique $Q$ of the graph. Thus, QSTAB($G$) may be seen as the set of vectors $x \in \mathbb{R}^{|V|}$ such that $x_i \geq 0$, and the QSTAB-inequalities associated with $G$ are satisfied.

Before we prove Theorem 1, note that the characterization of STAB($G$) was given in terms of its vertices and that of QSTAB($G$) was in terms of its hyperplanes. The following (known) link, Lemma 4, between these representations is key to the simplification.

Lemma 4 [39]. STAB($G$) is the convex hull of the integer solutions to the inequalities $x_i \geq 0$, and STAB-inequalities for $G$, where STAB inequalities for a graph $G = (V, E)$ are defined to be the set of inequalities given by $(x_i + x_j) \leq 1$ for every $(i, j) \in E$.

Proof of Theorem 1 We consider a five-cycle graph but our techniques readily generalize to the odd $n$-cycle case (unless stated otherwise). The QSTAB inequalities, together with the $x_i \geq 0$ condition, can be expressed as

$$0 \leq x_i \leq 1 \quad \text{for } i = \{1, 2, \ldots, 5\},$$

$$x_i + x_{i+1} \leq 1 \quad \text{for } i = \{1, 2, \ldots, 5\}$$

(3)

(4)

where $i + 1$ is modulo 5. Note the STAB inequalities, together with $x_i \geq 0$, turn out to be exactly the same as the aforesaid for the five-cycle graph. (The set STAB is a convex hull of integer solutions of STAB inequalities.) Each inequality is characterized by a hyperplane. The vertices must lie on the intersection of (at least) five distinct hyperplanes. From this, we can already see that the integer solutions of STAB inequalities and the QSTAB inequalities are the same. The FCNC inequality is one of the facets defining NC inequality. To see this, it suffices to observe that there are exactly five vertices of STAB, the corresponding behaviors of which saturate the said inequality (since the space is five dimensional) and remaining vertices satisfy the same inequality. For the five-cycle case, the remaining argument is trivial and we defer the proof of the $n$-cycle case to the end.
Suppose $x$ is a solution of QSTAB inequalities.
Then, the components of $x$ are either

all integers, all non-integers or mixed
(irrelevant) (unique) (no solution).

FIG. 1. To establish Theorem 1, it suffices to show that there is exactly one noninteger solution, i.e., a solution with at least one noninteger component, to QSTAB inequalities. We show that of the three kinds of solutions possible (as depicted above) the all-integer solution is irrelevant, the all-non-integer solution is unique, and no mixed solutions exist.

Note that, together with the aforementioned, if we can establish that there is only one noninteger solution of QSTAB inequalities then we have proven our result.

To this end, observe that there can only be the following three types of solutions: (1) all $x_i$ are integers, (2) none of the $x_i$ are integers, or (3) neither all $x_i$ are integers nor all $x_i$ are nonintegers (viz., at least one integer and at least one noninteger solution; see Fig. 1).

We are interested in the latter two cases. In case 2, we cannot use any of the QSTAB inequalities involving only one term [Eq. (3)]. This is because for a vertex we saturate five distinct inequalities. In this case, saturation of any of these inequalities will yield integer solutions which we are not considering. Hence, the only possibility is to use Eq. (4). Now we show that the solution is unique. Let $x_1 = q$ for any $0 < q < 1$. Saturating, we deduce $x_2 = 1 - q$, $x_3 = q$, $x_4 = 1 - q$, $x_5 = q$, and finally $x_1 = 1 - q$. This entails $x_1 = 1 - q = q$, which means $q = 1/2$ uniquely.

To complete the argument, we must show that there are no solutions in case 3. We already ruled out considering all five one term inequalities [Eq. (3)] as they yield integer solutions. Let us consider $k$ two term inequalities [Eq. (4)] and $m$ one term inequalities such that $m + k = 5$. The $m$ one term inequalities, when saturated (because we consider the intersection of hyperplanes to obtain the vertices), will force the corresponding $x_i$ to be integers. This means that there are at least $m$ integer $x_i$. To analyze further, we consider the following game. Consider the five-cycle graph (see Fig. 2). Select $m$ vertices of the graph (not to be confused with the vertices of QSTAB) and $k$ edges. The vertices correspond to the variables fixed by the one term inequalities (saturated, so equalities). The edges correspond to the two term inequalities (again, saturated so equalities). Two cases can arise in such an assignment. Either each of the $k$ edges is connected to one of the $m$ vertices (possibly via other edges, if not directly) or there is at least one edge which is not connected to any of the $m$ vertices (again, possibly via other edges, if not directly). These two cases are represented by the left and right graph in Fig. 2. Consider the second case. The disconnected edge (in the sense described earlier) will correspond to a two term equality involving two variables which have no other constraints. This means that the set of inequalities chosen does not uniquely determine a solution, i.e., at least one of the inequalities chosen is redundant. This case is therefore irrelevant. Consider the first case now. In this case, start with any one of the $m$ vertices. This corresponds to a one term equality which fixes the associated variable as an integer (as was noted earlier). Now the edge (if there is one) connected to this vertex directly will fix the value of the other vertex associated with the edge to be an integer. This reasoning can be repeatedly used to show that all the variables involved along the edges connected to the said initial vertex are integers. This can be repeated for every one of the $m$ vertices. This means that all variables are assigned integer values. We have reached a contradiction which means there are no solutions of the kind assumed by case 3.

We end by showing that the FCNC inequality is facet defining (in the exclusivity graph approach). All incidence vectors (we will restrict to the ones corresponding to the stable set of the $n$-cycle graph, for this proof) will always satisfy the FCNC inequality because the cardinality of the stable set is bounded by the independence number (see Definition 6) of the graph, which for our case is $(n - 1)/2$ [40–42]. We will now show that there are exactly $n$ vertices of STAB, i.e., incidence vectors which saturate the said inequality. To saturate, the incidence vector must have $(n - 1)/2$ components with entry one, and the remaining $(n + 1)/2$ components with entry zero. Note that each incidence vector satisfies the STAB inequalities, i.e., if a given component is one then its adjacent components are necessarily zero. One can convince themselves that any such vector, i.e., incidence vectors that saturate the FCNC inequality, must have two zeros adjacent (cyclically over $n$) while all other entries are alternatively one and zero. The total number of ways of placing two adjacent zeros, which is exactly $n$, then gives us the total number of incidence vectors which saturate the inequality, thereby proving that the FCNC inequality is indeed facet defining.

FIG. 2. There are two possible scenarios corresponding to the case where there is at least one integer and one noninteger solution (case 3 in the proof). The two term inequalities decide the values for two $x_i$s and have been represented as edges, and nodes (highlighted as small circles) have been used to denote the values determined by the one term inequalities. Depending on the way the combination of inequalities is selected, one gets either all $x_i$s as integers or a redundant set of inequalities leading to an undecidable value for $x_i$s.

IV. UNIQUENESS OF FUNDAMENTAL ANTICYCLIC INEQUALITIES

We show that the odd anticycle admits a unique inequality. This follows easily from the following known result. For any set of nonnegative vectors $X$, its antiblocker is defined as

$$\text{abl} X = \{ y \geq 0 : x, y \leq 1 \ \forall \ x \in X \}.$$
v \in X \) because every element \( y \in ab X \) will satisfy \( \{\theta_1 x_1 + \ldots + \theta_n x_n, y \leq 1 \} \) where \( \sum_{i=1}^{n} \theta_i = 1 \).

**Theorem 2.** Let \( G \) be odd \( n \)-cycle graph \( (n > 3) \) and consider the exclusivity graph scenario associated with \( G \). There is a unique facet-defining NC inequality, given by \( \sum_{i=1}^{n} p_i \leq 2 \).

**Proof.** We characterize QSTAB(\( G \)) and STAB(\( G \)) by using Lemma 5. Let \( \{v_1, v_2, \ldots, v_n\} \) denote the vertices of STAB(\( G \)). Each vertex \( v_l \) corresponds to a hyperplane constraining abl STAB(\( G \)) = QSTAB(\( G \)). From the proof of Theorem 1 we know that QSTAB(\( G \)) has exactly one more vertex, call it \( v_0 \). Corresponding to \( v_0 \) there will be exactly one extra hyperplane constraining abl QSTAB(\( G \)) = STAB(\( G \)) compared to those constraining abl STAB(\( G \)) = QSTAB(\( G \)). This hyperplane is precisely \( \sum_{i=1}^{n} p_i \leq 2 \) using \( v_0 = (1/2, 1/2, \ldots, 1/2) \) and the definition of the antiblocker. \( \square \)

**V. GEOMETRIC REPRESENTATION**

Having proved the uniqueness of the fundamental noncontextuality inequalities, we briefly discuss the geometry of the associated polytopes.

Fix an odd \( n > 3 \). Geometrically, the FCNC inequality corresponds to a unique hyperplane cutting through QSTAB which separates all E-principle behaviors (see Definition 31 in Appendix B) uniquely into two parts, namely, noncontextual and contextual. Naïvely one might imagine the QSTAB polytope and the FCNC inequality to geometrically be illustrated by the image on the left in Fig. 3. However, it is not too hard to show that all the facets of QSTAB are also the facets of STAB, which means the naïve understanding is flawed. The image on the right in Fig. 3 better illustrates the geometry of the two convex polytopes (STAB and QSTAB). Recall that STAB for a given graph is identical to the antiblocker of QSTAB for the complement graph and similarly QSTAB for a given graph is identical to the antiblocker of STAB for the complement graph (see Lemma 5). We thus recover essentially the same geometry for the anticyclic exclusivity scenario. In summary, given an \( n \)-cyclic (or anticyclic) exclusivity scenario, the associated fundamental NC inequality separates the corresponding E-principle behaviors into parts and uniquely singles out a vertex corresponding to the maximally contextual behavior.

**VI. ALL BELL INEQUALITIES FOR THE SIMPLEST EXCLUSIVITY GRAPH**

Any exclusivity scenario witnessing contextuality will have associated with it at least one fundamental NC inequality but there may be others; we give an example of this using the simplest Bell exclusivity scenario and report a Bell inequality in the process.

Given a list \( [L] \) of integers, a graph with \( n \) vertices where every \( i \)th vertex is connected to every other \( \ell \) mod \( n \)th vertex with \( \ell \in [L] \) is called a circulant graph \( C_n[L] \). The exclusivity graph corresponding to the measurement events for the CHSH inequality is a circulant graph. In particular, it is represented as the \( C_n[1, 4] \) graph (see Fig. 4). This is the simplest exclusivity graph which can lead to Bell inequalities if one analyzes the corresponding stable set polytope [36].

We ran numerical tests to characterize the convex polytope STAB(\( C_n[1, 4] \)) and found three types of nontrivial facets.

1. **Pentagonal inequalities:** There are eight such inequalities and they correspond to eight induced pentagons (five cycles). In the literature, these are known as pentagonal Bell inequalities [36]. Formally, the pentagonal inequalities are given by

\[
\sum_{i=k}^{(k+4) \text{mod} 8} p_i \leq 2
\]

for \( k \in \{1, 2, 3, 4, 5, 6, 7, 8\} \).

2. **Heptagonal inequalities:** There are eight such inequalities and correspond to induced subgraphs of \( C_n[1, 4] \) with any seven nodes. Formally, the heptagonal NC inequalities are given by

\[
\sum_{i=k}^{(k+6) \text{mod} 8} p_i \leq 3
\]

for \( k \in \{1, 2, 3, 4, 5, 6, 7, 8\} \).
(3) CHSH inequality: The CHSH inequality corresponds to the sum of the probabilities of all eight events and is formally given by

\[ \sum_{i=1}^{8} p_i \leq 3. \]

Note that neither the Heptagonal inequalities nor the CHSH inequality correspond to any odd cycle (or its complement).

VII. RELATIVE SIMPLIFICATION EXPLAINED

Where is this exponential simplification coming from and are we losing something in the process? To answer these questions we connect the inequality we obtained using the exclusivity graph approach to the one obtained using the compatibility hypergraph approach. We briefly introduce the compatibility hypergraph approach first (deferring a complete description to Appendix C).

A. Compatibility hypergraph approach: Overview

Instead of looking at the exclusivity of events, the compatibility hypergraph approach is based on, as the name suggests, the compatibility of measurements. The scenario here is defined by a hypergraph \( G \), where the vertices \( V(G) \) represent the measurements and the hyperedges \( C(G) \) of the graph represent the set of measurements which are compatible (see Definition 33 and 34 in Appendix C). The set of nodes or measurements in a hypergraph constitute a context (see Definition 20 in Appendix A). For each context \( C \in C \), a normalized probability vector \( p_C \) is defined which assigns probabilities to every possible joint outcome corresponding to the measurements belonging to that context (see Definition 20 in Appendix A). The behavior \( p \), in this formalism, is defined to be the concatenation of all these probability vectors (one vector for each context; see Definition 37 in Appendix C). Consider a fixed measurement. The behaviors which assign the same marginal probability distribution to the outcomes corresponding to this measurement, regardless of which probability vector (and hence context) was used to evaluate the marginal, are said to be nondisturbing behaviors. The set of such behaviors happens to be a convex polytope and is denoted by \( \mathcal{S}(G) \) (pronounced as \( X \); see Definition 40 in Appendix C). It is analogous (and reduces, under the appropriate restrictions) to the no-signaling distributions. Consider the set of behaviors which arise by associating a quantum observable with each measurement such that the compatibility requirements are satisfied. This set is called the set of quantum behaviors \( \mathcal{C}(G) \) and it happens to be convex. Finally, consider behaviors which admit a noncontextual completion (see Definition 22 in Appendix A). These behaviors can be equivalently characterized as behaviors which arise from a global joint distribution over all the observables such that the marginals yield the various probability vectors constituting the behavior (see Theorem 3 in Appendix C). These define the set of classical behaviors \( \mathcal{E}(G) \) which, like \( \mathcal{S}(G) \), happens to be a convex polytope. The facet-defining hyperplanes of \( \mathcal{E}(G) \) yield NC inequalities.

B. KCBS inequality

For concreteness, we consider the KCBS scenario in the compatibility hypergraph formalism. Let \( \{M_0, M_1, M_2, M_3, M_4\} \) be five dichotomic measurements, which form the set of vertices \( V(G) \), and let \( \{M_0, M_1\}, \{M_1, M_2\}, \{M_2, M_3\}, \{M_3, M_4\}, \{M_4, M_0\} \) denote the compatibility relations among them \( C(G) \), which define the (hyper)edges. This results in \( G \) being a five-cycle (hyper)graph [see Fig. 5 (top)]. It would be helpful to use two different conventions for labeling the outcomes of the measurements: binary \((0, 1)\) and signed \((-1, 1)\). In the binary convention, the probability of obtaining the outcome \((0, 0)\) upon the measurement of \( M_i \) and \( M_{i+1} \) is denoted by \( p_{i,i+1}(00) \). This is consistent with our notation for the probability vector \( p_C \) where \( C \) was the context, which in this case is denoted by \( i, i+1 \). The addition in the indices is modulo 5. Similarly in the signed convention, the outcome \((+1, +1)\) is denoted by \( p_{i,i+1}(++) \). Further, to represent the expectation value of the product of measurements \( M_i M_{i+1} \) we use the notation

\[ \langle i, i+1 \rangle = p_{i,i+1}(++) + p_{i,i+1}(-) - p_{i,i+1}(-+). \]

Note that we discuss expectation values only in the signed convention. By studying the polytope of classical behaviors \( \mathcal{E}(G) \) one can find the facet-defining hyperplanes which yield the NC inequalities. These turn out to be [35]

\[ \gamma_0(0, 1) + \gamma_1(1, 2) + \gamma_2(2, 3) + \gamma_3(3, 4) + \gamma_4(4, 0) \leq 3 \]
where $\gamma_i \in \{\pm 1\}$ such that the number of $\gamma$s with a negative sign is odd. To see how these relate to the unique KCBS inequality we obtain using the exclusivity graph approach, consider the possible events corresponding to a given context [the (hyper)edge]. All these events are mutually exclusive because they correspond to different outcomes of a given set of measurement [see Fig. 5 (bottom)]. This exclusivity is denoted by the ellipse in the graph. Now consider the measurement $M_0$. Let $(01|M_0M_1)$ denote the event where $M_0$ and $M_1$ are measured and the outcomes are zero and one, respectively (see Definitions 23 and 24 in Appendix B). Using this notation, observe that the events $(01|M_0M_1), (00|M_0M_1), (11|M_0M_1),$ and $(01|M_0M_0)$ are mutually exclusive. The first two are exclusive as they correspond to different outcomes of $M_1$, the second and third are exclusive because they correspond to different outcomes of $M_0$, and the third and fourth are exclusive because they correspond to different outcomes of $M_4$. A similar argument works for all the other pairs. This exclusivity is denoted by a straight line. One can verify that all the events on a given straight line are exclusive. If we let $e_i = (01|M_0M_{i+1})$ then a five-cycle exclusivity graph can be extracted from the aforementioned by defining the set of vertices to be the set of events $\{e_0, e_1, \ldots, e_5\}$ and the set of edges to be their exclusivity relations $\{[e_0, e_1], [e_1, e_2], \ldots, [e_5, e_0]\}$. We had denoted the probability of these events by $p_i$ to write the KCBS inequality, which in our current notation becomes

$$p_{0,1}(01) + p_{1,2}(01) + p_{2,3}(01) + p_{3,4}(01) + p_{4,0}(01) \leq 2. \tag{6}$$

Since the exclusivity graph approach does not require the explicit specification of the measurements which lead to exclusivity, merely their existence, it is able to extract the essential nature of the problem without creating redundancies caused by the labeling. If the complete exclusivity graph was used then the exclusivity graph formalism should yield effectively [43] the same NC inequalities as the compatibility hypergraph approach, recreating the said redundancies. We now explicitly combine the following two NC inequalities obtained using two different exclusivity graphs,

$$\sum_{i=0}^{4} p_{i,i+1}(01) \leq 2, \quad \sum_{i=0}^{4} p_{i,i+1}(10) \leq 2,$$

into an NC inequality obtained using the compatibility hypergraph formalism:

$$-\sum_{i=0}^{4} \langle i, i+1 \rangle \leq 3, \tag{7}$$

which corresponds to taking all $\gamma_i = -1$. Using $p_{i,i+1}(00) = p_{i,i+1}(++)$ (which holds by assumption; they were just different labels for the same outcomes) in Eq. (5), and probability conservation (probability vectors $p_C$ are normalized; in a given context, the probabilities sum to one), we deduce

$$\langle i, i+1 \rangle = 1 - 2[p_{i,i+1}(01) + p_{i,i+1}(10)]. \tag{8}$$

To obtain Eq. (7) we sum the two inequalities

$$\sum_{i=0}^{4} [p_{i,i+1}(01) + p_{i,i+1}(10)] \leq 4 \quad \iff \quad \sum_{i=0}^{4} (2p_{i,i+1}(01) + p_{i,i+1}(10) - 1) \leq 8 - 5 = 3.

While at first sight it might appear that the inequality obtained using the compatibility hypergraph is weaker as it is a linear combination of two exclusivity graph based inequalities, this conclusion is incorrect. This is because we can do better. We can obtain from a single compatibility hypergraph based inequality a corresponding exclusivity graph based inequality. We show this explicitly for $\sum_{i=0}^{4} p_{i,i+1}(01) \leq 2$. The marginal $p_i(1) = \sum_{j=0}^{4} p_{i,j}(1)$ for every $i, j$ belonging to a hyperedge. We start with noting that $p_i(10) + p_i(01) = p_i(1) - p_i(11) + p_{i,i+1}(01)$ but we can also write $p_i(1) = p_{i-1,1}(01) + p_{i-1,1}(11)$ (this is a consequence of the no-disturbance requirement). Consequently, using Eq. (8), we have

$$1 - \langle i, i+1 \rangle = 2[p_{i-1,i}(01) + p_{i-1,i}(11) - p_{i,i+1}(11) + p_{i,i+1}(01)] \Rightarrow 5 - \sum_{i=0}^{4} \langle i, i+1 \rangle = 4 \sum_{i=0}^{4} p_{i,i+1}(01)$$

and substituting Eq. (7) we obtain $4 \sum_{i=0}^{4} p_{i,i+1}(01) \leq 5 + 3$, which is just Eq. (6).

The NC inequalities for the compatibility hypergraph structure corresponding to anticycles are not known in the literature. A numerical investigation for the classical polytope for the same indicates the presence of exponentially many nontrivial NC inequalities. However, in the exclusivity graph approach, the nontrivial NC inequalities corresponding to anticycles are unique (see Theorem 2). The reason for the simplification is the same as that for the cyclic case.

**VIII. DISCUSSION AND CONCLUSION**

We showed that all fundamental NC inequalities are unique for their corresponding odd cycle (or anticycle) exclusivity scenario. This is an exponential simplification compared to the analogous compatibility hypergraph scenarios. Thereafter, we discussed the geometry of the associated polytopes. Every exclusivity scenario which exhibits contextuality must admit at least one fundamental NC inequality; however, a behavior might exhibit contextuality by violating some other NC inequality. We reported a Bell inequality in the process of exemplifying this phenomenon. Finally, we also analyzed the source of the simplification by relating the KCBS inequality in the exclusivity graph approach with the compatibility hypergraph approach.

We outline two directions for further exploration.

1. The first direction is the linking of our results to the resource theory of contextuality. The simplification for cyclic and anticyclic exclusivity scenarios indicates that there is a unique way to capture the amount of contextuality, to wit,
TABLE I. The anticycle case appears naturally in the exclusivity
graph approach due to the strong perfect graph theorem. We show
that both the cyclic and anticyclic case in this approach admit a
unique NC inequality while Araújo et al. [35] obtained exponentially
many NC inequalities in the compatibility hypergraph approach. The
characterization of the anticycle case for the latter approach is an
open problem.

| Cycles | This paper | Araújo et al. [35] |
|--------|------------|-------------------|
| Anticycles | This paper | ? |

the distance of the contextual behavior from the hyperplane
corresponding to fundamental NC inequalities. The strong
perfect graph theorem could potentially allow a generalization
to all exclusivity scenarios.

(2) The second direction is the exploration of anticycles
in the compatibility hypergraph approach. We characterized
the NC polytopes corresponding to the cycles and anticycles
in the exclusivity graph approach and compared the former
with its compatibility hypergraph analog. The characterization
of the anticycle case in the latter approach is still open (see
Table 1).

We note that self-testing using the NC inequality for the
cyclic case was shown in [15] and it was recently extended to
the anticycle case as well [44]. Our results, together with the
aforementioned works on self-testing, show that this charac-
terization is complete in the following sense: in the cyclic and
anticyclic case, there is a unique (fundamental) NC inequality
and this can be self-tested. Further, by the strong perfect
graph theorem, every exclusivity scenario which can exhibit
contextuality must contain at least one such fundamental NC
inequality.

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APPENDIX A: PROBABILISTIC MODELS: STATES
AND MEASUREMENTS

The results discussed here are based on the work of Amiral
and Cunha [32]. In any experimental scenario there are two
types of interventions possible, either preparation or opera-
tion. Preparation is used in the intuitive sense of the word, that
is, preparing the system in a given state, for instance using
a laser to initialize the state of an atom. More explicitly, we
make the following assumptions about the theory.

(1) Interventions are of two types: Preparation and opera-
tion.

(2) Experiments are reproducible: For each operation, there
may be several different outcomes, each occurring with a
well-defined probability for a given preparation.

**Definition 11. State.** Two preparations are defined to be
equivalent if they give the same probability distribution for
all available operations. We will refer to an equivalence class
of preparations as a state state.

**Definition 12. State space.** The set of all states is referred
to as the state space of the system.

**Remark 2.** The state space is convex.

**Definition 13. Pure states.** All extremal points of the state
space are defined to be pure states.

**Definition 14. Measurements.** Measurements are operations
with more than one outcome.

**Definition 15. Probabilistic model.** We call any mathemat-
ical description of a physical system which provides the
following a probabilistic model: (1) objects to represent [(a)
state, (b) operations, and (c) measurements] and (2) a rule
to calculate the probabilities of the possible outcomes of any
arbitrary measurement given any arbitrary state.

**Definition 16. Probability theory.** A probability theory is a
collection of probabilistic models.

**Definition 17. Outcome repeatable measurements.** A mea-
surement is defined to be an outcome repeatable measure-
ment if every time one performs this measurement on a system
and an outcome k is obtained a subsequent measurement
of j on the same system gives the outcome k again with
probability 1.

**Definition 18. \( p_j(k) \).** The probability of getting an outcome
k given that a measurement j has been performed will be
denoted by \( p_j(k) \).

All the measurements henceforth will be assumed to be
outcome repeatable.

**Definition 19. Compatible measurements, refinement, and
coarse graining.** A set of measurements \([j_1, j_2 \ldots j_n] \)
is compatible if there is another measurement j with outcomes
\([1, \ldots m] \) and functions \([f_1, f_2 \ldots f_m] \) such that the possible
outcomes of \( j_s \) are the same as \( \{f_1(1), f_2(2) \ldots f_m(m)\} \) for each
s and

\[
p_j(l) = \sum_{k \in f^{-1}(l)} p_j(k)
\]

where j is called a refinement of \([j_1, j_2 \ldots j_n] \) and each \( j_s \) is
called a coarse graining of j.

If a set of measurements is compatible it is called a set of
compatible measurements.

Completion of a probabilistic model

**Definition 20. Context.** A set of compatible measurements
is defined to be a context.

Our objective now is to construct a general mathematical
framework which can describe the completion of a probabilis-
tic model, i.e., give a model which is no longer probabilistic
but reduces to the same probabilistic model if certain variables
are ignored.
Definition 21. Completion. Consider a probabilistic model $P$ where $S$ represents the set of pure states and $X$ represents the set of measurements. The completion of this probabilistic model, denoted by $P'$, consists of a set of measurements $X'$, which are in one-to-one correspondence with $X$, and a set of pure states $S'$, which are in one-to-one correspondence with $\Lambda \times S$ for some set $\Lambda$. $P'$ must satisfy the following requirements. For all $\rho \in S$ and all contexts $c = \{j_1, j_2, \ldots, j_n\}$, $P'$ should specify a probability distribution over $\Lambda$ given by $p(\lambda)$ and a probability distribution $p^{(i,\rho,c)}_{ji} : \mathbb{R} \to \{0, 1\}$ for each $\lambda, \rho, c, j_k$ such that

\[ p^c(i_1, i_2, \ldots, i_m) = \sum_{\lambda \in \Lambda} p(\lambda) \prod_{k=1}^m p^{(i_k,\rho,c)}_{ji}(i_k) \]

where $p^c(i_1, i_2, \ldots, i_m)$ is the probability assigned by $P$ to the measurement of $j_1, j_2, \ldots, j_n$ (encoded in $c$) yielding the outcomes $i_1, i_2, \ldots, i_m$, respectively, for the state $\rho$.

Remark 4. We expect the completion $P'$ to specify $S'$ as $(\lambda, \rho)$ for all $\lambda \in \Lambda$ and $\rho \in S$. Let us assume for simplicity that $X' = X$. Now for every context $c = \{j_1, j_2, \ldots, j_n\}$ (i.e., set of compatible measurements from $X$) the completion $P'$ will predict with certainty the outcome of measuring any $j_k \in c$, for a given $(\lambda, \rho)$. This prediction is allowed to depend on the set $c$ itself to accommodate “contextual completions.” We will see later that noncontextual (and functionally consistent) completions contradict the predictions of quantum mechanics.

APPENDIX B: THE EXCLUSIVITY GRAPH APPROACH

1. Formalizing scenarios

Definition 22. Noncontextual completion. Let $c_1 = \{j_1, j_2, \ldots, j_n\}, c_2 = \{j_1', j_2', \ldots, j_n'\}$ be two contexts (note that $j_i$ and $j_i'$ may not be compatible for $i, k > 1$). A completion $P'$ of a probabilistic model $P$ is called noncontextual if $p^{(i,\rho,c_1)}_{ji} = p^{(i,\rho,c_2)}_{ji}$ for all contexts $c_1$ and $c_2$ of the aforesaid form.

Definition 24. Exclusive event. Two events $e_i$ and $e_j$ are defined to be exclusive if there exists a measurement $M$ such that $e_i$ and $e_j$ correspond to different outcomes of $M$, i.e., $e_i = (a_i|M)$ and $e_j = (a_j|M)$ such that $a_i \neq a_j$.

Definition 25. Exclusivity graph. For a family of events $\{e_1, e_2, \ldots, e_n\}$ we associate a simple undirected graph, $G := (V, E)$, with vertex set $V$ and edge set $E$ such that two vertices $i, j \in V$ share an edge if and only if $e_i$ and $e_j$ are exclusive events. $G$ is called an exclusivity graph.
2. Impossible completions: Linking geometry and quantum mechanics

Lemma 6 [4]. Let \( e_1, e_2 \ldots e_n \) be the (exclusive) events associated with an exclusivity graph \( G = (V,E) \). Then,
\[
B_{NC}(G) = \text{STAB}(G),
B_{O}(G) = \text{TH}(G),
B_{E}(G) = Q\text{STAB}(G).
\]

Corollary 1. For a given exclusivity graph \( G \) we have
\[
B_{NC}(G) \subseteq B_{O}(G) \subseteq B_{E}(G).
\]

APPENDIX C: COMPATIBILITY HYPERGRAPH APPROACH

1. Formalizing scenarios

Definition 33. Compatibility scenario. A compatibility scenario is specified by the tuple \( \Upsilon := (X, C, O) \) where the following hold.
1. \( O \) is a finite set.
2. \( X \) is a finite set of random variables from \( O \) to \( \mathcal{P}(O) \).
3. \( C \) is a collection of subsets of \( X \) such that their union is equal to \( X \) and the intersection of any two subsets is never equal to one of the subsets. Each of these subsets will be called a context.

It might be useful to keep the following in mind to get an intuitive understanding. The set \( X \) can be thought of as the measurements, the set \( O \) as the outcome of these measurements, and the set \( C \) as containing maximal contexts.

Definition 34. Compatibility hypergraph. The compatibility hypergraph corresponding to a scenario \( \Upsilon := (X, C, O) \) is a hypergraph the nodes of which are elements in \( X \) and hyperedges are contexts in \( C \). We denote the compatibility hypergraph for the scenario by \( \mathcal{H} = (X, C) \).

We use calligraphic letters to notationally distinguish the (hyper)graphs associated with the compatibility hypergraph from those associated with the exclusivity graph approach.

Definition 35. Compatibility graph. The compatibility graph for a scenario is the two-skeleton of the corresponding hypergraph, i.e., \( \mathcal{H} = (X, C) \). We will denote the compatibility graph by \( G = (X,E) \). Given two elements \( i \) and \( j \) in \( X \), they share an edge in \( E \) if and only if \( \{i , j \} \subseteq C \) for some context \( C \in C \).

Definition 36. Measurement event. A measurement event corresponds to a single run of an experiment where the measurements in a context \( C \in C \) are jointly performed with outcomes in \( O^C \).

Henceforth, for notational simplicity, we use \( O^C \) instead of \( O(\mathcal{H}) \).

Definition 37. Behavior and behavior vector. Given a scenario \( \Upsilon := (X, C, O) \), a behavior is a family of probability distributions defined over \( O^C \), defined as
\[
B = \left\{ p_C : O^C \rightarrow [0,1] \mid \sum_{s \in O^C} p_C(s) = 1, C \in \mathcal{C} \right\}.
\]

One can stack the \( p_C(s) \) for all \( s \in O^C \) and \( C \in \mathcal{C} \) to form a column vector of probabilities for a behavior \( B \). Such probability vectors are called behavior vectors.

Remark 8. The set of possible behaviors forms a polytope with \( \prod_{C \in \mathcal{C}} |O^C| \) nodes, where the extreme points (nodes) correspond to deterministic points, i.e., \( p_C(s) \) equal to zero or one. This can be deduced from the convexity of probability distributions.

Definition 38. Restriction map. Given a context \( C \in \mathcal{C} \) with outcomes in \( O^C \) and a set \( U \subset C \), a restriction map \( r^C_U \) is given by
\[
r^C_U : O^C \rightarrow O^U,
\]
\[
s = (a_i)_{i \in C} \mapsto s|_U = (a_i)_{i \in U}.
\]

Definition 39. Marginal distribution (for a context). The marginal distribution for a probability distribution over a context \( C \in \mathcal{C} \) corresponding to a set \( U \subset C \) is defined as
\[
p^C_U(s) = \sum_{r \in O^U \times \prod_{i \in \langle r \rangle \cap C} \mathbb{P}(C)} p_C(r),
\]
where \( r^C_U \) is a restriction map from \( O^C \) to \( O^U \).

Definition 40. Nondisturbing set, \( \mathcal{X}(\Upsilon) \). Given a scenario \( \Upsilon := (X, C, O) \), the set of behaviors is called a nondisturbing set, if for any given behavior and two different contexts \( C_1 \) and \( C_2 \) we have \( p^C_{1 \cap C_2} = p^C_{1 \cup C_2} \).

Definition 41. Global section. Given a scenario \( \Upsilon := (X, C, O) \), a global section for \( X \) is a probability distribution over \( O^X \), denoted by \( p_X : O^X \rightarrow [0,1] \).

Definition 42. Global section for a behavior. Given a scenario \( \Upsilon := (X, C, O) \) and a behavior \( B \), a global section \( p_X \) is called a global section for the behavior \( B \) if
\[
p^C_B = p_X \quad \forall \ C \in \mathcal{C}.
\]

Definition 43. Noncontextual behavior. A behavior which admits a global section is called a noncontextual behavior.

Remark 9. Noncontextual completions were defined independently.

Theorem 3. Abramsky and Brandenburger [2]. Given a scenario \( \Upsilon := (X, C, O) \) and a behavior \( B \in \mathcal{X}(\Upsilon) \), the aforementioned behavior \( B \) has a global section if and only if there exists a noncontextual completion recovering \( B \).

2. Probability distributions and physical theories

a. Classical realizations and noncontextuality

Remark 10. Given a scenario \( \Upsilon \), the set of classical behaviors will be denoted by \( \mathcal{C}(\Upsilon) \). Note that this notation is distinct from the \( C \) which was used to define the compatibility cover for the scenario \( \Upsilon \).

Claim 1. Given a scenario \( \Upsilon \) and a behavior \( B \), the following statements are equivalent.
1. \( B \) has a global section.
2. \( B \) is classical.

Remark 11. There exists a noncontextual completion recovering \( B \).

b. Quantum realization

Remark 12. The set of quantum behaviors will be denoted by \( \mathcal{Q}(\Upsilon) \).

Theorem 4. \( \mathcal{Q}(\Upsilon) \) is a convex set.
Remark 13. Restricting the dimension of realization in the quantum case yields a nonconvex set.

c. Noncontextuality inequalities

Remark 14. Given a scenario $\mathcal{Y}$ and a behavior $B$, how can one determine if there exists a noncontextual completion recovering $B$? This motivates us to define linear inequalities, the violation of which (for a behavior $B$) guarantees that $B$ is contextual, i.e., there is no noncontextual completion recovering $B$. It is important to note that the set of noncontextual behaviors forms a polytope, which means that characterization of the same can be given by intersection of finitely many hyperplanes and half spaces. NC inequalities for a scenario correspond to the facets of the classical (or noncontextual) polytope. Note that the set of nondisturbing behaviors $\mathcal{X}(\mathcal{Y})$ is also a polytope.

Given a polytope, the representation in terms of half spaces and hyperplanes is often called $H$ representation. The same polytope can also be described as a convex hull of finitely many vertices of the polytope, referred to as $V$ representation.

Definition 44. $\mathcal{C}$ inequality. The NC inequality is a linear inequality

$$S := \sum_{C \in \mathcal{C}, c \in C} \gamma_{c}(s)P_{C}(s) \leq b$$

where $\gamma_{C}(s)$ and $b$ are real numbers such that the inequality is satisfied for the behaviors in the noncontextual polytope $\mathcal{C}(\mathcal{Y})$. Often $b$ is called the noncontextual hidden variable bound because the noncontextual behaviors [the behaviors in $\mathcal{C}(\mathcal{Y})$] respect the bound.

Remark 15. There may exist behaviors in $\mathcal{X}(\mathcal{Y})$ which violate one or many NC inequalities for the scenario $\mathcal{Y}$. Such behaviors are often called contextual behaviors. An NC inequality is called tight if it exists noncontextual behavior saturating the inequality. Furthermore, it is called facet defining if it corresponds to one of the facets of the noncontextual polytope. Given a behavior $B$, its membership in $\mathcal{C}$ is equivalent to checking if all the facet defining NC inequalities are satisfied.

d. Geometry of the case $\mathcal{H} = \mathcal{G}$ (compatibility hypergraph = compatibility graph)

Remark 16. If we assume that every context has at most two measurements, then the compatibility hypergraph for the scenario is given by its two-skeleton. Furthermore, if we assume that every measurement has two outcomes, it leads to description of classical and no-disturbing sets as familiar polytopes from graph theory.

Description of the nondisturbing quantum and noncontextual behaviors.

Remark 17. The nondisturbing set $\mathcal{X}(\mathcal{Y})$ lies in $\mathbb{R}^{4|E|}$ because every edge (or equivalently context) corresponds to two binary measurements.

Definition 45. Notation. Fix $\{M_{i}, M_{j}\} \in \mathcal{C}$.

1. $p_{ij}(ab)$ is the probability of outcome $a$ and $b$ for the joint measurement of $i$ and $j$, respectively.
2. $p_{i}(a) = \sum_{b} p_{ij}(ab)$.
3. $p_{j}(b) = \sum_{a} p_{ij}(ab)$.

Claim 2. $p_{ij}(ab)$ can be determined from $p_{ij}(11)$ and $p_{i}(1)$ due to the constraints on nondisturbing behaviors.

Definition 46. $\phi : \mathbb{R}^{4|E|} \rightarrow \mathbb{R}^{\left|V(G)\right|+\left|E(G)\right|}$, $B \mapsto q = (q_{1}, q_{k})_{(i \in V(G), (i, j) \in E(G))}$ where $q_{i} = p_{i}(1)$ and $q_{ij} = p_{ij}(11)$.

Remark 18. To return from the $q$ space to the $B$ space, we use

$$p_{ij}(10) = q_{i} - q_{ij},$$
$$p_{ij}(01) = q_{j} - q_{ij},$$
$$p_{ij}(00) = 1 - q_{i} - q_{j} + q_{ij}.$$

Remark 19. The map is injective.

Definition 46. Correlation vector $v(S)$, correlation polytope. Correlation vector $v(S)$ is defined as follows. Given $S \subset V(G)$, $v(S) \in \mathbb{R}^{\left|V(G)\right|+\left|E\right|}$ is defined as

$$v(S)_{i} = \begin{cases} 1 & i \in S \\ 0 & \text{else} \end{cases} \forall i \in V,$$
$$v(S)_{ij} = \begin{cases} 1 & i, j \in S \\ 0 & \text{else} \end{cases} \forall \{i, j\} \in E.$$

The correlation polytope is defined to be the convex hull of all correlation vectors.

Theorem 5. $\phi[\mathcal{C}(\mathcal{Y})] = \text{COR}(G)$.

Definition 47. (For completeness) Rooted correlation semimetric polytope. $\text{RCMET}(G)$ of a graph $G$ is the set of vectors $q = (q_{1}, q_{kj}) \in \mathbb{R}^{\left|V(G)\right|+\left|E(G)\right|}$ such that

$$q_{ij} \geq 0,$$
$$q_{i} - q_{ij} \geq 0,$$
$$1 - q_{i} - q_{j} + q_{ij} \geq 0.$$

Theorem 6. $\phi[\mathcal{X}(G)] = \text{RCMET}(G)$.

e. Cut polytope

Definition 48. Cut vector, cut polytope. Given a graph $G$ and $c \in \{0, 1\}^{\left|V(G)\right|}$ the cut vector is $x(c) \in \mathbb{R}^{\left|E(G)\right|}$ such that

$$x(C)_{ij} = c_{i} \oplus c_{j}.$$

The cut-01 polytope, $\text{CUT}^{01}(G)$, is the convex hull of all cut vectors of $G$.

Definition 49. ±1 cut vectors. Given a graph $G$ and $c \in \{0, 1\}^{\left|V(G)\right|}$ the cut vectors are defined as $y(c) \in \mathbb{R}^{\left|E(G)\right|}$ such that

$$y(C)_{ij} = c_{i}c_{j}.$$

The cut ±1 polytope, $\text{CUT}^{\pm1}(G)$, is the convex hull of all ±1 cut vectors of $G$.

Definition 50. Suspension graph ($V(G)$). The suspension graph $V(G)$ of $G$ is the graph with vertex set $V(G) \cup \{e\}$ and edge set $E(G) \cup \{(e, i), i \in V(G)\}$.

Remark 20. $V(G)$ is the graph obtained by adding an extra vertex and joining all the vertices of $G$ to it.

Theorem 7. $\text{CUT}^{01}(V(G)) = \text{COR}(G)$, where $\psi : \mathbb{R}^{\left|V(G)\right|+\left|E(G)\right|} \rightarrow \mathbb{R}^{\left|V(G)\right|+\left|E(G)\right|}$ is defined by

$$q \mapsto x.$$
where

\[ x_{ij} = 1 - q_i - q_j + 2q_{ij}, \quad \{i, j\} \in E(G), \]
\[ x_{ei} = q_i, \quad i \in V(G). \]

**Remark 21.** \( x_{ei} = \langle j \rangle = q_j = p_j(1) \) where \( \langle j \rangle \) is the expectation value corresponding to the observable \( j \). Similarly, \( x_{ij} = p_{ij}(00) + p_{ij}(11) \) is the probability of getting the same outcome.

**Claim 3.** \( \text{CUT}^{01} \) and \( \text{CUT}^{01} \) are related by a bijective linear map \( \alpha : \text{CUT}^{01}(G) \rightarrow \text{CUT}^{01}(G) \) defined by

\[ x \mapsto y \]

where

\[ y_{ij} = 2x_{ij} - 1. \]

**Remark 22.** If we relabel the outcomes as \( 0 \rightarrow +1 \) and \( 1 \rightarrow -1 \) then we can write

\[ y_{ei} = -\langle i \rangle = -p_{j}(1) - p_{j}(-1) \]

and

\[ y_{ij} = -(ij) \]

\[ = -[p_{ij}(1, 1) + p_{ij}(-1, -1) - p_{ij}(-1, 1) - p_{ij}(1, -1)]. \]

Now we discuss an example.

**Example 2. Bell scenario.** The Bell scenario corresponds to the case where context is generated via spacelike separation of the involved parties. We explain the defining components of the Bell scenario \( \Upsilon := (X, C, O) \) hereafter. We assume the number of parties to be \( n \). The set \( X \) consists of various disjoint subsets \( X_1, X_2, \ldots, X_n \). The subset \( X_i \) consists of measurement operators for party \( i \). All contexts \( C \in \mathcal{C} \) are of the form \( \mathcal{C} = \{M_1, M_2, \ldots, M_n\} \) where \( M_i \in X_i \). The compatibility graph for the scenario is a complete \( n \)-partite graph. The Bell scenario corresponding to \( n \) parties with \( m \) measurements per party where each measurement has \( o \) outcomes is often denoted as \( (n, m, o) \).

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