Using concave optimization methods for inexact quadratic programming problems with an application to waste management

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Abstract

Quadratic programming is potentially capable of strategic decision making in real world problems. However, practical problems rarely conform to crisp parameters, and hence the prospects of these problems with inexact parameters are inevitably higher. The existing studies regarding public welfare schemes/organizations reveal that their objectives end up as minimization of cost functions and are governed by linear or concave quadratic programming problems. The present study proposes a method that can be applied to concave type quadratic objective function subject to linear constraints with inexact parameters. A comparison is also drawn with existing methods to establish its simplicity and efficiency. Further, a numerical example is illustrated, and finally, a waste management problem is formulated and solved using the proposed method.

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1 Introduction

Waste management (WM) problem is paramount, in particular, both to the public and environment. Human rate of consumption and consequently production of various wastes has far outrun nature’s capacity to rejuvenate. Improper handling of the problem can cause considerable damage and can even trigger a catastrophic situation. WM affects our daily activities and is of prime concern to the environmentalists, administrators, and industry. The WM issue has drawn huge attention from researchers with focus on different aspects of this problem. Mathematicians, in particular, are trying to fit a suitable programming model with imprecise parameters, which can in general be applied to WM. Primarily, linear programming (LP) model with interval parameters happens to be the first choice due to its simplicity. These models have been used to find solution of many real life problems viz. WM, ground water allocation, and flood control.

There have been many studies related to WM on various directions. Chang and Lu [1] worked on creating WM planning alternatives in view of fuzzy environmental resources.
Further, Sakai and Hiraoka [2] investigated WM solution by focusing on recycling of incinerator residue. Thereafter, Cai et al. [3] proposed a model with a combination of interval, fuzzy, and stochastic approach for WM. Next, Fan et al. [4] introduced a solution method for fuzzy LP to reduce computational requirements in the WM problem. Sun et al. [5] worked on the WM problem using a constrained quadratic model and provided an effective tool for reflecting system cost variations. Further, Tan et al. [6] identified optimal plans in fuzzy environment using a two-step solution method. On a different note, Lu et al. [7] and He et al. [8] focussed on controlling gas emission in WM using dynamic optimization and mixed integer programming approach, respectively. Thereafter, Astrup et al. [9] investigated the incinerator performance linked to emissions and residues. Li and Chen [10] suggested fuzzy stochastic interval parameters in LP model for WM. Next, Fan and Huang [11] proposed a two-step method for WM to avoid loss of decision related information, while Li and Huang [12] demonstrated a robust quadratic programming model to handle long term WM planning and decisions. Later, Ionescu et al. [13] stressed on an integrated WM model using pretreatment and revenue generating processes. Rada et al. [14] proposed an integrated solution for WM by maintaining a balance between energy and environment. Next, Xu et al. [15] presented a model to deal with possibilistic or probabilistic uncertainties and tackle complexities derived from capacity-expansion issues.

In reference to the various methodologies proposed in the literature, Sugimoto et al. [16] introduced a parallel relaxation method for quadratic programming problems. Later, Van Thorai [17] presented a finite branch and bound algorithm as an optimization technique for a quadratic programming problem. Further, Gao et al. [18] proposed a rectangle branch and reduce approach for faster convergence of nonconvex quadratic programming algorithm. Next, Liu and Wang [19] proposed a numerical solution method for interval quadratic programming. Fan et al. [20] illustrated an effective robust method for interval LP in uncertain environmental decision making. Thereafter, Fan et al. [21] described generalized fuzzy method for uncertainties in the form of fuzzy sets and developed a method using $\alpha$-cut and membership functions.

On a different note, Atalay and Apaydin [22] suggested a method based on gamma distribution for stochastic programming model while a method for nonlinear optimization using quadratic programming was proposed by Zhu et al. [23]. Next, an algorithm for generalized vector quasi-complementarity problems with an application to traffic network problems was proposed by Van Hung et al. [24]. Recently, Van Hung et al. [25] discussed convergence of solution sets for fuzzy optimization problems, and then a new class of generalized multiobjective games with fuzzy mappings was proposed by Van Hung et al. [26]. Further, Van Hung and Keller [27] improvised the traffic network problem in [24] to a traffic network under uncertainty by establishing convergence of fuzzy vector quasi-optimization problems.

In the present day scenario of a continuously changing and competitive market, the decision making has to be handled with utmost care when crisp parameters fail to provide a satisfactory solution. The volatile data in practical problems can be conveniently accommodated using interval parameters in comparison to all other types of imprecise parameters. Using interval parameters in an LP model, Li et al. [28] studied flood control planning. Further, Li and Huang [29] introduced a stochastic quadratic programming method to handle uncertainties and nonlinearities in water management. Later, Fan et al. [30] explained a generalized fuzzy LP model for air quality management. Further, Jin et al. [31]
investigated allocation of water for irrigation purposes. Miao et al. [32] worked on an interval fuzzy programming method for water resource planning. An inexact water quality management model supporting economic and environmental management was proposed by Zhang et al. [33]. Thereafter, Huang et al. [34] studied controlling of noise in urban localities with the help of interval data. Though the LP models with interval parameters take care of the uncertainty part, governing all costs by a constant function may not tackle every practical situation.

As an application to WM under imprecise environment with linearly decreasing costs, Chen and Huang [35] initiated the application of quadratic programming problem using inexact parameters. The solution of resulting concave quadratic programming problem was found by applying derivative algorithm. The WM problem was further investigated by Kong et al. [36] and solved using duality approach in an effort to give an alternate method. However, duality results are applicable under certain conditions without which the results are misleading and do not comply with the characteristics of duality theorem. The present article not only provides an alternate strategy to tackle the WM problem, but also highlights and rectifies the discrepancy in Kong et al. [36].

The proposed algorithm is based on a two-level interval programming procedure which ultimately yields an interval solution of the given problem. The remaining part of the article is organized as follows. Section 2 reviews some basic definitions. Section 3 develops a methodology to solve inexact concave quadratic programming problems. Section 4 illustrates an example to display the steps involved in the proposed methodology. Section 5 deals with the application of inexact quadratic programming in the WM problem using the proposed approach. Next, a counterexample is constructed to describe the discrepancy in Kong et al. [36]. Further, computational efficiency of the proposed algorithm and some limitations of the inexact quadratic programming model are also discussed. Finally, conclusions with some future directions are presented in Sect. 6.

2 Basic definitions

Throughout the paper, $\mathbb{R}$ denotes the set of real numbers. Some of the definitions related to interval parameters are given below.

**Definition 2.1** ([36]) An interval number $A^\pm$ is defined as an interval with known upper and lower bounds, but unknown distribution information and is defined as follows:

$$A^\pm = [A^-, A^+] = \{a \in \mathbb{R} | A^- \leq a \leq A^+\},$$

where $A^-$ and $A^+$ are the lower and upper bounds of $A^\pm$, respectively.

**Arithmetic operations on intervals**

Suppose that $A^\pm$ and $B^\pm$ are two intervals, then for $k_1, k_2 \in \mathbb{R}$,

$$k_1[A^-, A^+] + k_2[B^-, B^+] = \begin{cases} [k_1A^- + k_1B^-, k_1A^+ + k_2B^+] & \text{if } k_1 \geq 0, k_2 \geq 0, \\ [k_1A^- + k_2B^+, k_1A^+ + k_2B^+] & \text{if } k_1 \geq 0, k_2 \leq 0, \\ [k_1A^+ + k_2B^-, k_1A^+ + k_2B^+] & \text{if } k_1 \leq 0, k_2 \geq 0, \\ [k_1A^+ + k_2A^+, k_1A^- + k_2B^+] & \text{if } k_1 \leq 0, k_2 \leq 0. \end{cases}$$
Definition 2.2 ([36]) The ordering on two intervals \( A^\pm \) and \( B^\pm \) is defined as follows:

\[
A^\pm \leq B^\pm \quad \text{iff} \quad A^- \leq B^- \quad \text{and} \quad A^+ \leq B^+,
\]

\[
A^\pm < B^\pm \quad \text{iff} \quad A^- < B^- \quad \text{and} \quad A^+ < B^+.
\]

3 The inexact quadratic programming problem

To capture real world uncertainties, usual precise parameters do not serve the purpose. Among various methods to handle imprecise data, use of inexact or interval parameters is one of the simplest ways to handle realistic problems. An inexact quadratic programming problem can be formulated in the following manner:

\[
\begin{align*}
\text{Minimize} & \quad Z = \sum_{j=1}^{n} \left( C_j^+ x_j + D_j^+ x_j^2 \right) \\
\text{subject to} & \quad \sum_{j=1}^{n} A_{ij}^\pm x_j \leq B_i^\pm, \quad i = 1, 2, \ldots, m, x_j \geq 0, j = 1, 2, \ldots, n,
\end{align*}
\]

where, for all \( i \) and \( j \), \( C_j^\pm = [C_j^-, C_j^+] \), \( D_j^\pm = [D_j^-, D_j^+] \), \( A_{ij}^\pm = [A_{ij}^-, A_{ij}^+] \), and \( B_i^\pm = [B_i^-, B_i^+] \). Further, it is assumed that \( D_j^+ \leq 0 \) for all \( j \).

As \( Z \) fluctuates due to inexact parameters in an objective function as well as in constraints, it lies between two values \( Z^- \) and \( Z^+ \) (say), that is, \( Z = [Z^-, Z^+] \). Now, to obtain these lower and the upper bounds of (IQP), the problem is further divided into the following two sub-problems:

\[
\begin{align*}
Z^- &= \min_S \left( \min_x \left( \sum_{j=1}^{n} (C_j x_j + D_j x_j^2) \right) \right) \\
\text{subject to} & \quad \sum_{j=1}^{n} A_{ij} x_j \leq B_i, x_j \geq 0, i = 1, 2, \ldots, m, j = 1, 2, \ldots, n
\end{align*}
\]

and

\[
\begin{align*}
Z^+ &= \max_S \left( \min_x \left( \sum_{j=1}^{n} (C_j x_j + D_j x_j^2) \right) \right) \\
\text{subject to} & \quad \sum_{j=1}^{n} A_{ij} x_j \leq B_i, x_j \geq 0, i = 1, 2, \ldots, m, j = 1, 2, \ldots, n
\end{align*}
\]

where \( S = \{ (C_j, D_j, A_{ij}, B_i) \} \) \( C_j \in [C_j^-, C_j^+] \), \( D_j \in [D_j^-, D_j^+] \), \( A_{ij} \in [A_{ij}^-, A_{ij}^+] \), \( B_i \in [B_i^-, B_i^+] \), \( \forall i \) and \( j \).

Now, both the problems above involve bi-level optimization criteria (max-min type for \( Z^+ \) and min-min type for \( Z^- \)) with varying feasible region. Hence, to compute the bounds, it is necessary to convert these problems into single level optimization models with fixed feasible region. Before discussing the equivalent problems to evaluate these bounds, first we shall compute the smallest and largest feasible region of the problem (IQP). Consider
the following extreme cases of the feasible region:

\[ S_1 = \left\{ x \in \mathbb{R}^n \mid \sum_{j=1}^{n} A_{ij}^+ x_j \leq B_i^+, x_j \geq 0, \forall i, j \right\}, \]

\[ S_2 = \left\{ x \in \mathbb{R}^n \mid \sum_{j=1}^{n} A_{ij}^- x_j \leq B_i^-, x_j \geq 0, \forall i, j \right\}, \]

\[ S_3 = \left\{ x \in \mathbb{R}^n \mid \sum_{j=1}^{n} A_{ij}^+ x_j \leq B_i^+, x_j \geq 0, \forall i, j \right\}, \]

\[ S_4 = \left\{ x \in \mathbb{R}^n \mid \sum_{j=1}^{n} A_{ij}^- x_j \leq B_i^-, x_j \geq 0, \forall i, j \right\}. \]

Let \( x \in S_1 \) and as \( A_{ij}^- \leq A_{ij}^+ \), \( B_i^- \leq B_i^+ \), \( \forall i \) and \( j \), therefore

\[
\sum_{j=1}^{n} A_{ij}^- x_j \leq \sum_{j=1}^{n} A_{ij}^+ x_j \leq B_i^+, \quad \forall i.
\]

Hence, by the above inequality \( x \in S_2 \). Similarly, \( x \in S_2 \) implies \( x \in S_2 \) and \( x \in S_4 \) yields \( x \in S_2 \). This implies \( S_1 \cup S_2 \cup S_3 \cup S_4 = S_2 \). Thus, the largest feasible region of (IQP) is \( S_2 \).

Further, \( x \in S_1 \) yields \( x \in S_2 \), \( x \in S_3 \), and \( x \in S_4 \). That is, \( S_1 \subseteq S_2, S_1 \subseteq S_3, S_1 \subseteq S_4 \) and hence \( S_1 \cap S_2 \cap S_3 \cap S_4 = S_1 \). This proves that the smallest feasible region of (IQP) is \( S_1 \).

Next, we shall discuss the problems to evaluate \( Z^- \) and \( Z^+ \).

### 3.1 Model to compute \( Z^- \)

Observe the following facts about the problem (LIQP):

1. It has the same nature of optimization (min-min type).
2. \( C_j^- \leq C_j, D_j^- \leq D_j, x_j \geq 0 \) for all \( j \) and hence

\[
\sum_{j=1}^{n} (C_j^- x_j + D_j x_j^2) \leq \sum_{j=1}^{n} (C_j x_j + D_j x_j^2).
\]

3. Among all the feasible regions, the largest feasible region \( S_2 \) defined by the constraints of the problem (LIQP) yields the best minimum value of the objective function,

it finally reduces to

\[
Z^- = \min_x \left( \sum_{j=1}^{n} (C_j^- x_j + D_j x_j^2) \right) \tag{LIQPP}
\]

subject to \( \sum_{j=1}^{n} A_{ij}^- x_j \leq B_i^+, \quad i = 1, 2, \ldots, m, x_j \geq 0, j = 1, 2, \ldots, n. \)

### 3.2 Model to compute \( Z^+ \)

To solve (UIQP), which involves different nature of optimization viz. max-min type, we proceed as follows:
As \( C_j \leq C_j^* \), \( D_j \leq D_j^* \), and \( x_j \geq 0 \) for all \( j \), therefore
\[
\sum_{j=1}^{n} (C_j x_j + D_j x_j^2) \leq \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2).
\]

This further implies
\[
\min_x \left( \sum_{j=1}^{n} (C_j x_j + D_j x_j^2) \right) \leq \min_x \left( \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2) \right).
\]

Hence
\[
\max_S \left( \min_x \left( \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2) \right) \right) = \max_{S'} \left( \min_x \left( \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2) \right) \right),
\]

where \( S' = \{ (A_{ij}, B_i) \mid A_{ij} \in [A_{ij}^-, A_{ij}^+], B_i \in [B_i^-, B_i^+] \} \). Therefore, the problem (UIQP) is equivalent to
\[
Z^* = \max_{S'} \left( \min_x \left( \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2) \right) \right)
\]
subject to \( \sum_{j=1}^{n} A_{ij} x_j \leq B_i, \quad x_j \geq 0, \quad i = 1, 2, \ldots, m, j = 1, 2, \ldots, n. \)

Further, from the varying feasible regions, following best/ worst value criterion, we get
\[
\min_{S_{1,x}} \left( \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2) \right) \leq \min_{S',x} \left( \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2) \right) \leq \min_{S_{1,x}} \left( \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2) \right)
\]
This yields
\[
\max_{S'} \left( \min_x \left( \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2) \right) \right) = \min_{S_{1,x}} \left( \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2) \right).
\]
That is, out of all the possible, feasible regions corresponding to \( S' \), the best value of \( Z^* \) will be obtained over the smallest region \( S_1 \) since it gives the worst value of
\[
\min_x \left( \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2) \right).
\]

So, the problem of finding \( Z^* \) is finally reduced to
\[
Z^* = \min_x \left( \sum_{j=1}^{n} (C_j^* x_j + D_j^* x_j^2) \right) \quad \text{(UIQPP)}
\]
subject to \( \sum_{j=1}^{n} A_{ij} x_j \leq B_i, \quad i = 1, 2, \ldots, m, x_j \geq 0, j = 1, 2, \ldots, n. \)
Remark 1

1. It is to be noted that, for all \( j \), \( D^-_j \leq 0 \) and \( D^+_j \leq 0 \), which yield that the Hessian matrices of the problems (LIQPP) and (UIQPP) are negative semi-definite, and hence the problems are concave optimization models.

2. Due to the concave nature of the objective functions of (LIQPP) and (UIQPP), duality theory is also not applicable as shown in Kong et al. [36].

The following theorem guarantees the global optimal solution if the objective function (minimization type) is concave.

**Theorem 1** ([37]) Consider the following problem:

Minimize \( Z(x) \) over \( S = \{ x : G_i(x) \leq 0, i = 1, 2, \ldots, m \} \),

where \( Z \) is a concave function defined on \( \mathbb{R}^n \), \( G_i \) (\( i = 1, 2, \ldots, m \)) are convex functions defined on \( \mathbb{R}^n \) whose gradients are continuous, and \( S \) is compact such that a point in the strict interior of the feasible region exists. Then there exists an extreme point \( \hat{x} \) of the convex compact set \( S \) which globally minimizes the problem (P).

Flow chart representation of the proposed algorithm can be viewed in Fig. 1.

4 Numerical example

Minimize \( f = [2, 5]x_1 + [3, 4]x_2 - [3, 6]x_1^2 - [4, 5]x_2^2 \)
subject to \( [2, 3]x_1 + [3, 4]x_2 \leq [10, 12], \)
\( [2, 3]x_1 + [1, 2]x_2 \geq [4, 6], \)
\( x_1, x_2 \geq 0. \)

**Solution:** Using LIQPP and UIQPP models to find lower and upper bound, respectively, the two subproblems are:

**Lower bound:**
\[
\begin{align*}
\text{Minimize } & \quad 2x_1 + 3x_2 - 6x_1^2 - 5x_2^2 \\
\text{subject to } & \quad 2x_1 + 3x_2 \leq 12, \\
& \quad 3x_1 + 2x_2 \geq 4, \\
& \quad x_1, x_2 \geq 0,
\end{align*}
\]

**Upper bound:**
\[
\begin{align*}
\text{Minimize } & \quad 5x_1 + 4x_2 - 3x_1^2 - 4x_2^2 \\
\text{subject to } & \quad 3x_1 + 4x_2 \leq 10, \\
& \quad 2x_1 + x_2 \geq 6, \\
& \quad x_1, x_2 \geq 0.
\end{align*}
\]

Noting that the objective functions of the above problems are negative semi-definite, the model happens to be concave optimization. Using MAPLE-12, we get \( f^- = -204 \) (at \( x_1 = 6, x_2 = 0 \)) and \( f^+ = -16.67 \) (at \( x_1 = 3.33, x_2 = 0 \)). Hence, \( f = [-204, -16.67] \).

### 5 Application to waste management

The proposed method is applied to a WM study (Fig. 2) on three localities and two WM facilities viz. the landfill and the incinerator. Total time span is fifteen years, which is split into three time spans of five years each. Due to impreciseness in all of collection and transportation costs, landfill and incinerator operating costs, residue disposal cost, and earnings from generated energy sales, an interval quadratic programming model is most appropriate. The capacity of the landfill is approximately \( 3.75 \times 10^6 \) to \( 4.00 \times 10^6 \) tonnes and the incinerator capacity is 500 to 650 tonnes per day. It is observed that the incinerator leaves residues of nearly 30% of the incoming waste. The problem is to allocate appropriate waste flows from the localities to the suitable facility while minimizing the aggregate cost.

Taking into account the transportation costs, running costs of facilities, and earnings from energy generated out of the incinerator as some of the factors while determining the problem, the IQP problem to be investigated is formulated as follows:

Minimize \( f \) = the sum of

1. Collection and transportation cost: \( \sum_{i=1}^{3} \sum_{j=1}^{3} \sum_{k=1}^{3} L_k TR_{ijk}^x x_{ijk} \),
2. Running cost of the landfill/ incinerator: \( \sum_{i=1}^{3} \sum_{j=1}^{3} \sum_{k=1}^{3} L_k OP_{ijk} x_{ijk} \),
3. Cost of disposal of incinerator residue in the landfill: \( \sum_{j=1}^{3} \sum_{k=1}^{3} L_k FE(FT_k^x + OP_{2k}^x) x_{2jk} \),
4. Earnings from sale of energy: \( -\sum_{j=1}^{3} \sum_{k=1}^{3} L_k RE_k x_{2jk} \)

subject to

1. Landfill capacity constraint: \( \sum_{j=1}^{3} \sum_{k=1}^{3} L_k (x_{ijk} + FEx_{2jk}) \leq TL^x, \)
2. Incinerator capacity constraint: \( \sum_{j=1}^{3} x_{2jk} \leq TE_k^x, \forall k, \)
3. Waste disposal demand constraint: \( \sum_{i=1}^{2} x_{ijk} \geq WG_{jk}^x, \forall i, j, k, \)
4. Nonnegativity constraint: \( x_{ijk} \geq 0, \forall i, j, k. \)
where:

$f$ is the aggregate cost of waste disposal ($\$$),

$i = 1, 2$ stands for the facility; (1 for the landfill and 2 for the incinerator),

$j = 1, 2, 3$ stands for the locality,

$k = 1, 2, 3$ stands for the time span,

$L_k$ is length of the time span $k$ (=1825 days for all $k$),
Table 3 Waste transportation cost (in dollars)

|                | Time span (k = 1)                            | Time span (k = 2)                            | Time span (k = 3)                            |
|----------------|---------------------------------------------|---------------------------------------------|---------------------------------------------|
| $TR_{11k}^\pm$ | $[-0.0161X + 15.08, -0.0121X + 19.50]$     | $[-0.0181X + 16.02, -0.0139X + 21.30]$     | $[-0.0165X + 17.65, -0.0139X + 23.45]$     |
| $TR_{12k}^\pm$ | $[-0.0142X + 12.63, -0.0105X + 16.89]$     | $[-0.0164X + 13.92, -0.0116X + 18.58]$     | $[-0.0169X + 15.31, -0.0127X + 20.35]$     |
| $TR_{13k}^\pm$ | $[-0.0170X + 15.30, -0.0130X + 20.50]$     | $[-0.0188X + 16.84, -0.0142X + 22.54]$     | $[-0.0209X + 18.45, -0.0158X + 24.70]$     |
| $TR_{21k}^\pm$ | $[-0.0130X + 11.55, -0.0098X + 15.40]$     | $[-0.0144X + 12.75, -0.0109X + 16.97]$     | $[-0.0155X + 14.01, -0.0116X + 18.66]$     |
| $TR_{22k}^\pm$ | $[-0.0136X + 12.17, -0.0103X + 16.15]$     | $[-0.0149X + 13.40, -0.0113X + 17.75]$     | $[-0.0164X + 14.75, -0.0124X + 19.55]$     |
| $TR_{23k}^\pm$ | $[-0.0120X + 10.05, -0.0090X + 14.02]$     | $[-0.0130X + 11.65, -0.0098X + 15.52]$     | $[-0.0143X + 12.83, -0.0108X + 17.06]$     |

Table 4 Transportation cost from the incinerator to the landfill (in dollars)

|                | Time span (k = 1)                            | Time span (k = 2)                            | Time span (k = 3)                            |
|----------------|---------------------------------------------|---------------------------------------------|---------------------------------------------|
| $FT_{k}^\pm$  | $[-0.0065X + 5.70, -0.0047X + 7.65]$        | $[-0.0070X + 6.27, -0.0055X + 8.38]$        | $[-0.0076X + 6.90, -0.0058X + 9.25]$        |

Table 5 Earnings (dollars) and capacity of the incinerator (in tonnes/day)

|                | Time span (k = 1)                            | Time span (k = 2)                            | Time span (k = 3)                            |
|----------------|---------------------------------------------|---------------------------------------------|---------------------------------------------|
| $RE_k^\pm$    | $[15, 20]$                                  | $[15, 20]$                                  | $[15, 20]$                                  |
| $TE_k^\pm$    | $[500, 650]$                                | $[500, 650]$                                | $[500, 650]$                                |

$x_{ijk}$ is the waste transferred to facility $i$ from locality $j$ during time span $k$ (tonnes/day), $TR_{ijk}^\pm$ is collection and transportation cost from locality $j$ to facility $i$ during time span $k$ ($/tonne)$, $OP_{ik}^\pm$ is running cost of facility $i$ during time span $k$ ($/tonne)$, $FE$ is residue flow rate from the incinerator to the landfill (30%), $FT_{k}^\pm$ is transportation cost for residue from the incinerator to the landfill during time span $k$ ($/tonne)$, $RE_k^\pm$ is earnings from the incinerator in time span $k$ ($/tonne)$, $TL_k^\pm$ is capacity of the landfill ($3.75 \times 10^6 – 4.00 \times 10^6$ tonnes), $TE_k^\pm$ is capacity of the incinerator in time span $k$ (tonnes/day), $WG_{jk}^\pm$ is waste generation rate in locality $j$ during time span $k$ (tonnes/day).

The waste collection and transportation cost are governed by market trend of decreasing linear functions, and in particular for locality 1, represented by Figs. 3–6. Similarly, transportation cost for residue from the incinerator to the landfill is also observed following decreasing linear function trend.

Using data in the given model from Tables 1–5, the inexact quadratic programming problem becomes:

$$
\text{minimize } f = 1825 \left([50.08, 57]x_{111} + [74.51, 84.595]x_{211} + [47.63, 54.39]x_{121} + [75.13, 85.345]x_{221} + [50.30, 58]x_{131} + [73.01, 83.215]x_{231} + [66.02, 74.3]x_{112} + [92.13, 104.584]x_{212} + [63.92, 71.58]x_{122} + [92.78, 105.364]x_{222} + [66.84, 75.54]x_{132} + [91.03, 103.13]x_{232}\right)
$$
Figure 3 Lower bound of waste collection and transfer cost from locality 1 to the landfill

Figure 4 Upper bound of waste collection and transfer cost from locality 1 to the landfill

\[ + [81.65, 88.45]x_{113} + [104.63, 120.19]x_{213} + [79.31, 85.35]x_{123} \\
+ [105.37, 121.075]x_{223} + [82.45, 89.7]x_{133} + [103.45, 118.585]x_{233} \\
+ [-0.0161, -0.0121]x_{211} + [-0.01495, -0.01121]x_{311} \\
+ [-0.0142, -0.0105]x_{221} + [-0.01555, -0.01441]x_{321} \\
+ [-0.0170, -0.0130]x_{331} + [-0.01395, -0.01041]x_{231} \\
+ [-0.0181, -0.0139]x_{112} + [-0.0165, -0.01255]x_{212} \\
+ [-0.0146, -0.0116]x_{122} + [-0.017, -0.01295]x_{222} \\
+ [-0.0188, -0.0142]x_{132} + [-0.0151, -0.01145]x_{232} \\
+ [-0.0195, -0.0146]x_{113} + [-0.01778, -0.01354]x_{213} \\
+ [-0.0169, -0.0127]x_{123} + [-0.01868, -0.01414]x_{233} \]
\[ + [-0.0209, -0.0158]x_{133}^2 + [-0.01658, -0.01254]x_{233}^2 \]

subject to

\[ 1825(x_{111} + x_{121} + x_{131} + x_{112} + x_{132} + x_{113} + x_{123} + x_{133} + 0.3(x_{211} + x_{221} + x_{231} + x_{212} + x_{222} + x_{232} + x_{213} + x_{223} + x_{233})) \]
\[ \leq [3.75, 4.00] \times 10^6, \]
\[ x_{211} + x_{221} + x_{231} \leq [500, 650], \]
\[ x_{212} + x_{222} + x_{232} \leq [500, 650], \]
\[ x_{213} + x_{223} + x_{233} \leq [500, 650], \]
\[ x_{111} + x_{211} \geq [250, 350], \]
\[ x_{121} + x_{221} \geq [150, 250], \]
\[ x_{131} + x_{231} \geq [250, 350], \\
x_{112} + x_{212} \geq [300, 400], \\
x_{122} + x_{222} \geq [185, 265], \\
x_{132} + x_{232} \geq [250, 350], \\
x_{113} + x_{213} \geq [350, 440], \\
x_{123} + x_{223} \geq [215, 285], \\
x_{133} + x_{233} \geq [300, 400], \\
x_{ijk} \geq 0 \quad \forall i, j, k. \]

Now, using the model (LIQPP), we have

\[
\begin{align*}
\min_{x} & \quad 1825(50.08x_{111} + 74.51x_{211} + 47.63x_{121} + 75.13x_{221} + 50.30x_{131} \\
& + 73.01x_{231} + 66.02x_{112} + 92.131x_{212} + 92.781x_{222} + 66.84x_{132} \\
& + 91.031x_{232} + 81.65x_{113} + 104.63x_{213} + 79.31x_{123} + 105.37x_{223} + 82.45x_{133} \\
& + 103.45x_{233} - 0.0161x_{111}^2 - 0.01495x_{211}^2 - 0.0142x_{121}^2 - 0.01555x_{221}^2 \\
& - 0.0170x_{131}^2 - 0.01395x_{231}^2 - 0.0181x_{112}^2 - 0.0165x_{212}^2 - 0.0146x_{122}^2 - 0.017x_{222}^2 \\
& - 0.0188x_{132}^2 - 0.0151x_{232}^2 - 0.0195x_{113}^2 - 0.01778x_{213}^2 - 0.0169x_{123}^2 - 0.01868x_{223}^2 \\
& - 0.0209x_{133}^2 - 0.01658x_{233}^2)) \\
\text{subject to} & \quad 1825(x_{111} + x_{112} + x_{131} + x_{112} + x_{122} + x_{132} + x_{113} + x_{123} + x_{133} \\
& + 0.3(x_{211} + x_{221} + x_{231} + x_{212} + x_{222} + x_{232} + x_{213} + x_{223} + x_{233})) \\
& \leq 4.00 \times 10^6, \\
x_{211} + x_{221} + x_{231} \leq 650, \\
x_{212} + x_{222} + x_{232} \leq 650, \\
x_{213} + x_{223} + x_{233} \leq 650, \\
x_{111} + x_{211} \geq 250, \\
x_{121} + x_{221} \geq 150, \\
x_{131} + x_{231} \geq 250, \\
x_{112} + x_{212} \geq 300, \\
x_{122} + x_{222} \geq 185, \\
x_{132} + x_{232} \geq 250, \\
x_{113} + x_{213} \geq 350, \\
x_{123} + x_{223} \geq 215, \\
x_{133} + x_{233} \geq 250, \\
x_{112} + x_{122} + x_{132} \geq [185, 265], \\
x_{113} + x_{123} + x_{133} \geq [350, 440], \\
x_{211} + x_{221} + x_{231} \geq [250, 350], \\
x_{212} + x_{222} + x_{232} \geq [300, 400], \\
x_{213} + x_{223} + x_{233} \geq [185, 265].
\]
Table 6  Optimal values of variables for lower bound $f^-$

| Time span (k) | Facility type (i) | Locality (j) | Optimum waste allocation (tonnes) |
|---------------|------------------|--------------|-----------------------------------|
| 1             | 1                | 1            | $x_{111} = 250$                   |
|               |                  | 2            | $x_{121} = 150$                   |
|               |                  | 3            | $x_{131} = 250$                   |
| 2             | 1                | 1            | $x_{211} = 0$                     |
|               |                  | 2            | $x_{221} = 0$                     |
|               |                  | 3            | $x_{231} = 0$                     |
| 2             | 1                | 1            | $x_{112} = 300$                   |
|               |                  | 2            | $x_{122} = 102$                   |
|               |                  | 3            | $x_{132} = 250$                   |
| 3             | 1                | 1            | $x_{113} = 350$                   |
|               |                  | 2            | $x_{123} = 215$                   |
|               |                  | 3            | $x_{133} = 300$                   |
|               | 1                | 1            | $x_{211} = 0$                     |
|               |                  | 2            | $x_{221} = 0$                     |
|               |                  | 3            | $x_{231} = 0$                     |

Lower bound $f^-$ $\geq 260,670,295$

$x_{133} + x_{233} \geq 300,$

$x_{ijk} \geq 0 \quad \forall i,j,k.$

Solving using MAPLE-12, $f^- = 260,670,295$ at the values given in Table 6.

Further, using the model (UIQPP) for upper bound, we obtain

$$f^+ = \min_{x} 1825(57x_{111} + 84.595x_{211} + 54.39x_{121} + 85.345x_{221} + 58x_{131} + 83.215x_{231} + 74.3x_{112} + 105.584x_{212} + 71.58x_{122} + 105.364x_{222} + 75.54x_{132} + 103.134x_{232} + 88.45x_{113} + 120.185x_{213} + 85.35x_{123} + 121.075x_{223} + 89.7x_{133} + 118.585x_{233} - 0.0121x_{111}^2 - 0.01121x_{211}^2 - 0.0105x_{121}^2 - 0.01441x_{221}^2 - 0.0130x_{131}^2 - 0.01041x_{231}^2 - 0.0139x_{112}^2 - 0.01255x_{212}^2 - 0.0116x_{122}^2 - 0.01295x_{222}^2 - 0.0142x_{113}^2 - 0.01145x_{213}^2 - 0.0146x_{123}^2 - 0.01354x_{223}^2 - 0.0127x_{133}^2 - 0.01414x_{233}^2 - 0.0158x_{133}^2 - 0.01254x_{233}^2)$$

subject to

$$1825(x_{111} + x_{121} + x_{131} + x_{112} + x_{122} + x_{132} + x_{113} + x_{123} + x_{133} + x_{211} + x_{221} + x_{231} + x_{212} + x_{222} + x_{232} + x_{213} + x_{223} + x_{233}) \leq 3.75 \times 10^6,$$

$x_{211} + x_{221} + x_{231} \leq 500,$

$x_{212} + x_{222} + x_{232} \leq 500,$

$x_{213} + x_{223} + x_{233} \leq 500,$

$x_{111} + x_{211} \geq 350,$
Table 7 Optimal values of variables for upper bound $f^+$

| Time span (k) | Facility type (i) | Locality (j) | Optimum waste allocation (tonnes) |
|---------------|-------------------|--------------|----------------------------------|
| 1             | 1                 | 1            | $x_{111} = 0$                    |
|               |                   | 2            | $x_{211} = 250$                  |
|               |                   | 3            | $x_{311} = 200$                  |
| 2             | 1                 | 1            | $x_{112} = 0$                    |
|               |                   | 2            | $x_{122} = 165$                  |
|               |                   | 3            | $x_{222} = 400$                  |
|               |                   | 2            | $x_{232} = 0$                    |
| 3             | 1                 | 1            | $x_{113} = 0$                    |
|               |                   | 2            | $x_{123} = 285$                  |
|               |                   | 3            | $x_{133} = 361$                  |
|               |                   | 2            | $x_{213} = 440$                  |
|               |                   | 3            | $x_{233} = 39$                   |

Upper bound $f^+ = $ 473,418,686

$x_{121} + x_{221} \geq 250,$
$x_{131} + x_{231} \geq 350,$
$x_{112} + x_{212} \geq 400,$
$x_{122} + x_{222} \geq 265,$
$x_{132} + x_{232} \geq 350,$
$x_{113} + x_{213} \geq 440,$
$x_{123} + x_{223} \geq 285,$
$x_{133} + x_{233} \geq 400,$
$x_{ijk} \geq 0 \quad \forall i,j,k.$

It yields $f^+ = $ 473,418,686 (using MAPLE-12). The values at which the function attains the upper limit are given in Table 7. The graphical interpretation of the results can be seen in Figs. 7–8. Here, LF stands for the landfill and IC stands for the incinerator.

**Remark 2** The objective functions of the given WM problem dealing with $f^-$ and $f^+$ are of minimization type with Hessian matrices as diagonal of order 18 × 18 having all negative entries. This shows that the problems are having concave type objective functions subject to linear equations. Further, if we apply the duality principle for the concave minimization type objective function subject to linear constraints, the optimal values of the two problems may not be the same. For instance, consider the following example:

Min $f(x) = -x^2 - y^2 + x + y$  \hspace{1cm} (E1)
subject to $0 \leq x \leq 1, 0 \leq y \leq 1.$

It is to be noticed that the objective function in (E1) is concave subject to linear inequalities. Directly solving the problem, we get $x = 1, y = 1,$ and $f(x) = 0$ as the optimal solution.
However, on applying duality, we obtain

\[
\text{Max } f(x) = x^2 + y^2 - \lambda_1 - \lambda_2 \tag{E2}
\]

subject to

\[-2x + \lambda_1 - \delta_1 = -1,\]

\[-2y + \lambda_2 - \delta_2 = -1,\]

\[x \geq 0, y \geq 0, \lambda_1 \geq 0, \lambda_2 \geq 0, \delta_1 \geq 0, \delta_2 \geq 0,\]

which after solving yields \(x = 0.5, y = 0.5\), and \(f(x) = 0.5\). This implies that the optimal values of (E1) and (E2) are not the same. It shows that duality is not applicable in (E1) as it is not a convex optimization problem (see p. 232, Bazaraa et al. [38]).
Based on the above discussion, it can be concluded that the implementation of the duality based algorithm by Kong et al. [36] to a WM problem governed by an inexact concave quadratic programming is erroneous. The present article not only develops a method to solve IQP with concave type objective function, but also removes a deficiency in the execution of the duality algorithm to the WM problem in Kong et al. [36].

**Computational efficiency of the proposed algorithm versus existing methods**

In the literature, a vast number of WM problems follow a linear model due to the assumption of constant costs along with imprecise parameters. However, the concave optimization models dealing with WM in [12, 35] are tackled by derivative algorithm. Proposed algorithm scores over the existing algorithms in the following sense:

- The derivative algorithm introduced by Chen and Huang [35] involves $2n$ variables and $2(m + n)$ constraints, whereas the proposed algorithm uses only $n$ variables and $m + n$ constraints, thus leading to a more computationally efficient algorithm.
- The procedure for derivative algorithm in [12, 35] is more cumbersome for application purposes in comparison to proposed algorithm which requires to identify only sense of the objective function and signs of squared terms.
- The duality based algorithm proposed by Kong et al. [36] for convex type quadratic programming problems involves $2n + m$ variables and $4n + m$ constraints which are also significantly higher in comparison to $n$ variables and $m + n$ constraints in the proposed algorithm.

**6 Conclusions**

The problems of WM are commonly governed by either LP problems or concave quadratic programming problems, depending upon constant or decreasing linear cost functions, respectively. The present article provides a solution methodology to tackle concave quadratic programming problems having inexact parameters. It is easy to apply and more computationally efficient in comparison to the derivative algorithm used in Li and Huang [12], Chen and Huang [35], and duality approach in Kong et al. [36]. Further, the discrepancy in the implementation of the algorithm to the application in the WM problem by Kong et al. [36] is also removed. In future, it will be interesting to investigate the algorithm for general model of quadratic/ nonlinear programming problem with inexact parameters and decision variables. Further, parabolic cost functions may be considered while dealing with application in WM.

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**Competing interests**

The authors declare that they have no competing interests.

**Authors’ contributions**

All authors jointly worked on the results, and they read and approved the final manuscript.
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