Modeling Attrition in Recommender Systems with Departing Bandits

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Abstract

Traditionally, when recommender systems are formalized as multi-armed bandits, the policy of the recommender system influences the rewards accrued, but not the length of interaction. However, in real-world systems, dissatisfied users may depart (and never come back). In this work, we propose a novel multi-armed bandit setup that captures such policy-dependent horizons. Our setup consists of a finite set of user types, and multiple arms with Bernoulli payoffs. Each (user type, arm) tuple corresponds to an (unknown) reward probability. Each user’s type is initially unknown and can only be inferred through their response to recommendations. Moreover, if a user is dissatisfied with their recommendation, they might depart the system. We first address the case where all users share the same type, demonstrating that a recent UCB-based algorithm is optimal. We then move forward to the more challenging case, where users are divided among two types. While naive approaches cannot handle this setting, we provide an efficient learning algorithm that achieves $\tilde{O}(\sqrt{T})$ regret, where $T$ is the number of users.

1 Introduction

At the heart of online services spanning such diverse industries as media consumption, dating, financial products, and more, recommendation systems (RSs) drive personalized experiences by making curation decisions informed by each user’s past history of interactions. While in practice, these systems employ diverse statistical heuristics, much of our theoretical understanding of them comes via stylized formulations within the multi-armed bandits (MABs) framework. While MABs abstract away from many aspects of real-world systems they allow us to extract crisp insights by formalizing fundamental tradeoffs, such as that between exploration and exploitation that all RSs must face (Joseph et al. 2016; Liu and Ho 2018; Patil et al. 2020; Ron, Ben-Porat, and Shalit 2021). As applied to RSs, exploitation consists of continuing to recommend items (or categories of items) that have been observed to yield high rewards in the past, while exploration consists of recommending items (or categories of items) about which the RS is uncertain but that could potentially yield even higher rewards.

In traditional formalizations of RSs as MABs, the recommender’s decisions affect only the rewards obtained. However, real-life recommenders face a dynamic that potentially alters the exploration-exploitation tradeoff: Dissatisfied users have the option to depart the system, never to return. Thus, recommendations in the service of exploration not only impact instantaneous rewards but also risk driving away users and therefore can influence long-term cumulative rewards by shortening trajectories of interactions.

In this work, we propose departing bandits which augment conventional MABs by incorporating these policy-dependent horizons. To motivate our setup, we consider the following example: An RS for recommending blog articles must choose at each time among two categories of articles, e.g., economics and sports. Upon a user’s arrival, the RS recommends articles sequentially. After each recommendation, the user decides whether to “click” the article and continue to the next recommendation, or to “not click” and may leave the system. Crucially, the user interacts with the system for a random number of rounds. The user’s departure probability depends on their satisfaction from the recommended item, which in turn depends on the user’s unknown type. A user’s type encodes their preferences (hence the probability of clicking) on the two topics (economics and sports).

When model parameters are given, in contrast to traditional MABs where the optimal policy is to play the best fixed arm, departing bandits require more careful analysis to derive an optimal planning strategy. Such planning is a local problem, in the sense that it is solved for each user. Since the user type is never known explicitly (the recommender must update its beliefs over the user types after each interaction), finding an optimal recommendation policy requires solving a specific partially observable MDP (POMDP) where the user type constitutes the (unobserved) state (more details in Section 5.1). When the model parameters are unknown, we deal with a learning problem that is global, in the sense that the recommender (learner) is learning for a stream of users instead of a particular user.

We begin with a formal definition of departing bandits in Section 2 and demonstrate that any fixed-arm policy is prone to suffer linear regret. In Section 3, we establish the UCB-based learning framework used in later sections. We instantiate this framework with a single user type in Section 4, where we show that it achieves $\tilde{O}(\sqrt{T})$ regret for $T$ being the number of users. We then move to the more challenging case with two user types and two recommenda-
tion categories in Section 5. To analyze the planning problem, we effectively reduce the search space for the optimal policy by using a closed-form of the expected return of any recommender policy. These results suggest an algorithm that achieves $O(\sqrt{T})$ regret in this setting. Finally, we show an efficient optimal planning algorithm for multiple user types and two recommendation categories (Appendix A) and describe a scheme to construct semi-synthetic problem instances for this setting using real-world datasets (Appendix B).

1.1 Related Work

MABs have been studied extensively by the online learning community (Cesa-Bianchi and Lugosi 2006; Bubeck, Cesa-Bianchi et al. 2012). The contextual bandit literature augments the MAB setup with context-dependent rewards (Abbasi-Yadkori, Pál, and Szepesvári 2011; Slivkins 2019; Mahadik et al. 2020; Korda, Szörényi, and Li 2016; Lattimore and Szepesvári 2020). In contextual bandits, the learner observes a context before they make a decision, and the reward depends on the context. Another line of related work considers the dynamics that emerge when users act strategically (Kremer, Mansour, and Perry 2014). Further works (Appendix B).

More broadly, our RS learning problem falls under the domain of reinforcement learning (RL). Existing RL literature that considers departing users in RSs include Zhao et al. (2020b); Lu and Yang (2016); Zhao et al. (2020a). While Zhao et al. (2020b) handle users of a single type that depart the RS within a bounded number of interactions, our work deals with multiple user types. In contrast to Zhao et al. (2020b), we consider an online setting and provide regret guarantees that do not require bounded horizon. Finally, Lu and Yang (2016) use POMDPs to model user departure and focus on approximating the value function. They conduct an experimental analysis on historical data, while we devise an online learning algorithm with theoretical guarantees.

2 Departing Bandits

We propose a new online problem, called departing bandits, where the goal is to find the optimal recommendation algorithm for users of (unknown) types, and where the length of the interactions depends on the algorithm itself. Formally, the departing bandits problem is defined by a tuple $([M], [K], q, P, \Lambda)$, where $M$ is the number of user types, $K$ is the number of categories, $q \in [0,1]^M$ specifies a prior distribution over types, and $P \in (0,1)^{K \times M}$ and $\Lambda \in (0,1)^{K \times M}$ are the click-probability and the departure-probability matrices, respectively.

There are $T$ users who arrive sequentially at the RS. At every episode, a new user $t \in [T]$ arrives with a type $type(t)$. We let $q$ denote the prior distribution over the user types, i.e., $type(t) \sim q$. Each user of type $x$ clicks on a recommended category $a$ with probability $P_{a,x}$. In other words, each click follows a Bernoulli distribution with parameter $P_{a,x}$. Whenever the user clicks, she stays for another iteration, and when the user does not click (no-click), she departs with probability $\Lambda_{a,x}$ (and stays with probability $1 - \Lambda_{a,x}$). Each user $t$ interacts with the RS (the learner) until she departs.

We proceed to describe the user-RS interaction protocol. In every iteration $j$ of user $t$, the learner recommends a category $a \in [K]$ to user $t$. The user clicks on it with probability $P_{a,type(t)}$. If the user clicks, the learner receives a reward of $r_{i,j}(a) = 1$ if the user does not click, the learner receives no reward (i.e., $r_{i,j}(a) = 0$), and user $t$ departs with probability $\Lambda_{a,type(t)}$. We assume that the learner knows the value of a constant $\epsilon > 0$ such that $\max_{a,x} P_{a,x} \leq 1 - \epsilon$ (i.e., $\epsilon$ does not depend on $T$). When user $t$ departs, she does not interact with the learner anymore (and the learner moves on to the next user $t + 1$). For convenience, the departing bandits problem protocol is summarized in Algorithm 1.1

Having described the protocol, we move on to the goals of the learner. Without loss of generality, we assume that the online learner’s recommendations are made based on a

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1 We denote by $[n]$ the set $\{1, \ldots, n\}$.

2 We formalize the reward as is standard in the online learning literature, from the perspective of the learner. However, defining the reward from the user perspective by, e.g., considering her utility as the number of clicks she gives or the number of articles she reads induces the same model.
The performance of the learner is compared to that of the best policy, formally defined by the \textit{regret} for \(T\) episodes,

\[
R_T = T \cdot \mathbb{E}[V^{\pi^*}] - \sum_{t=1}^{T} V^{\pi_t} .
\]

The learner’s goal is to minimize the expected regret \(\mathbb{E}[R_T]\).

**Algorithm 1: The Departing Bandits Protocol**

\textbf{Input}: number of types \(M\), number of categories \(K\), and number of users (episodes) \(T\)

\textbf{Hidden Parameters}: types prior \(q\), click-probability \(P\), and departure probability \(\Lambda\)

1: for episode \(t \leftarrow 1, \ldots, T\) do
2: a new user with type \(\text{type}(t) \sim q\) arrives
3: \(j \leftarrow 1, \text{depart} \leftarrow \text{false}\)
4: while \text{depart} is \text{false} do
5: the learner picks a category \(a \in [K]\)
6: with probability \(P_{a,x}\), user \(t\) clicks on \(a\) and \(r_{t,j}(a) \leftarrow 1\); otherwise, \(r_{t,j}(a) \leftarrow 0\)
7: if \(r_{t,j}(a) = 0\) then
8: with probability \(\Lambda_{a,x}\): \text{depart} \leftarrow \text{true} and user \(t\) departs
9: the learner observes \(\sum_{j=1}^{N^{\pi}(t)} r_{t,j}(\pi_{t,j})\)
10: if \text{depart} is \text{false} then
11: \(j \leftarrow j + 1\)

\[\sum_{t=1}^{T} V^{\pi_t} = \sum_{t=1}^{T} \sum_{j=1}^{N^{\pi}(t)} r_{t,j}(\pi_{t,j}).\]

The performance of the learner is compared to that of the best policy, formally defined by the \textit{regret} for \(T\) episodes,

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R_T = T \cdot \mathbb{E}[V^{\pi^*}] - \sum_{t=1}^{T} V^{\pi_t} .
\]

\[
\text{Table 1: The departing bandits instance in Section 2.1}
\]

| Category 1 | Type \(x\) | Type \(y\) |
|------------|-------------|-------------|
| \(P_{1,x} = 0.3\) | \(P_{1,y} = 0.28\) |
| \(P_{2,x} = 0.4\) | \(P_{2,y} = 0.39\) |

\[
\text{Prior } q_x = 0.4, \quad q_y = 0.6
\]

\[
\text{2.1 Example}
\]

The motivation for the following example is two-fold. First, to get the reader acquainted with our notations; and second, to show why fixed-arm policies are inferior in our setting.

Consider a problem instance with two user types \((M = 2)\), which we call \(x\) and \(y\) for convenience. There are two categories \((K = 2)\), and given no-click the departure is deterministic, i.e., \(\Lambda_{a,x} = 1\) for every category \(a \in [K]\) and type \(x \in [M]\). That is, every user leaves immediately if she does not click. Furthermore, let the click-probability \(P\) matrix and the user type prior distribution \(q\) be as in Table 1.

Looking at \(P\) and \(q\), we see that Category 1 is better for Type \(x\), while Category 2 is better for Type \(y\). Notice that without any additional information, a user is more likely to be type \(y\). Given the prior distribution, recommending Category 1 in the first round yields an expected reward of \(q_x \cdot P_{1,x} + q_y \cdot P_{1,y} = 0.368\). Similarly, recommending Category 2 in the first round results in an expected reward of 0.394. Consequently, if we recommend myopically, i.e., without considering the user type, always recommending Category 2 is better than always recommending Category 1.

Let \(\pi^a\) denote the fixed-arm policy that always selects a single category \(a\). Using the tools we derive in Section 5 and in particular Theorem 5.3, we can compute the expected returns of \(\pi^1\) and \(\pi^2\), \(\mathbb{E}[V^{\pi^1}]\) and \(\mathbb{E}[V^{\pi^2}]\). Additionally, using results from Section 5.2, we can show that the optimal policy for the planning task, \(\pi^*\), recommends Category 2 until iteration 7, and then recommends Category 1 for the rest of the iterations until the user departs.

Using simple calculations, we see that \(\mathbb{E}[V^{\pi^*}] - \mathbb{E}[V^{\pi^1}] > 0.0169\) and \(\mathbb{E}[V^{\pi^*}] - \mathbb{E}[V^{\pi^2}] > 1.22 \times 10^{-5}\); hence, the expected return of the optimal policy is greater than the returns of both fixed-arm policies by a constant. As a result, if the learner only uses fixed-arm policies (\(\pi^a\) for every \(a \in [K]\)), she suffers linear expected regret, i.e., \(\mathbb{E}[R_T] = T \cdot \mathbb{E}[V^{\pi^*}] - \sum_{t=1}^{T} \mathbb{E}[V^{\pi_t}] = \Omega(T)\).

\[
\text{3 UCB Policy for Sub-exponential Returns}
\]

In this section, we introduce the learning framework used in the paper and provide a general regret guarantee for it.

In standard MAB problems, at each \(t \in [T]\), the learner picks a single arm and receives a single sub-Gaussian reward. In contrast, in departing bandits, at each \(t \in [T]\) the learner receives a return \(V^\pi\), which is the cumulative reward of that policy. The return \(V^\pi\) depends on the policy \(\pi\) not only through the obtained rewards at each iteration

\[\text{We limit the discussion to deterministic policies solely; this is w.l.o.g. (see Subsection 5.1 for further details).}\]
There exist \( \tilde{\tau}, \tilde{\eta} \). Furthermore, as we have shown in Section 2.1, in terms of the number of users, which suggests considering a more expressive set of policies. This in turn yields another disadvantage for using MAB algorithms for departing bandits, as their regret is linear in the number of arms (categories) \( K \).

As we show later in Sections 4 and 5, for some natural instances of the departing bandits problem, the return from each user is sub-exponential (Definition 3.1). Algorithm 2, {\large \text{Algorithm 2}} below, receives a set of policies \( \Pi \) as input, along with other parameters that we describe shortly. The algorithm is a restatement of the UCB-Hybrid Algorithm from Jia, Shi, and Shen (2021), with two modifications: (1) The input includes a set of policies rather than a set of actions/categories, and accordingly, the confidence bound updates are based on return samples (denoted by \( V^n \)) rather than reward samples. (2) There are two global parameters (\( \tilde{\tau} \) and \( \tilde{\eta} \)) instead of two local parameters per action. If the return from each policy in \( \Pi \) is sub-exponential, Algorithm 2 not only handles sub-exponential returns, but also comes with the following guarantee: Its expected value is close to the value of the best policy in \( \Pi \).

### 3.1 Sub-exponential Returns
For convenience, we state here the definition of sub-exponential random variables (Eldar and Kutyniok, 2012).

**Definition 3.1.** We say that a random variable \( X \) is sub-exponential with parameters \( (\tilde{\tau}, b) \) if for every \( \gamma \) such that \( |\gamma| < 1/b \),

\[
\mathbb{E} [ \exp(\gamma(X - \mathbb{E}[X]))] \leq \exp(\frac{\gamma^2 \tilde{\tau}^2}{2}).
\]

In addition, for every \( (\tilde{\tau}^2, b) \)-sub-exponential random variables, there exist constants \( C_1, C_2 > 0 \) such that the above is equivalent to each of the following properties:
1. **Tails:** \( \forall v \geq 0 : \text{Pr}[|X| > v] \leq \exp(1 - \tilde{\tau}^2 v) \).
2. **Moments:** \( \forall p \geq 1 : \left( \mathbb{E}[|X|^p] \right)^{1/p} \leq C_2 p \).

Let \( \Pi \) be a set of policies with the following property: There exist \( \tilde{\tau}, \tilde{\eta} \) such that the return of every policy \( \pi \in \Pi \) is \( (\tilde{\tau}^2, b) \)-sub-exponential with \( \tilde{\tau} \geq \tilde{\tau} \) and \( \tilde{\eta} \geq \tilde{\eta} \). The following Algorithm 2 receives as input a set of policies \( \Pi \) with the associated parameters, \( \tilde{\tau} \) and \( \tilde{\eta} \). Similarly to the UCB algorithm, it maintains an upper confidence bound \( U \) for each policy, and balances between exploration and exploitation. Theorem 3.2 below shows that Algorithm 2 always gets a value similar to that of the best policy in \( \Pi \) up to an additive factor of \( \tilde{\tilde{O}} \left( \sqrt{|\Pi|T} + |\Pi| \right) \). The theorem follows directly from Theorem 3 from Jia, Shi, and Shen (2021) by having policies as arms and returns as rewards.

**Theorem 3.2.** Let \( \Pi \) be a set of policies with the associated parameters \( \tilde{\tau}, \tilde{\eta} \). Let \( \pi_1, \ldots, \pi_T \) be the policies Algorithm 2 selects. It holds that

\[
\mathbb{E} \left[ \max_{\pi \in \Pi} \left( V^{\pi} - \sum_{t=1}^{T} V^{\pi_t} \right) \right] = O(\sqrt{|\Pi|T \log T} + |\Pi| \log T).
\]

There are two challenges in leveraging Theorem 3.2. The first challenge is crucial: Notice that Theorem 3.2 does not imply that Algorithm 2 has a low regret; its only guarantee is w.r.t. the policies in \( \Pi \) received as an input. As the number of policies is infinite, our success will depend on our ability to characterize a “good” set of policies \( \Pi \). The second challenge is technical: Even if we find such \( \Pi \), we still need to characterize the associated \( \tilde{\tau} \) and \( \tilde{\eta} \). This is precisely what we do in Section 4 and 5.

### 4 Single User Type
In this section, we focus on the special case of a single user type, i.e., \( M = 1 \). For notational convenience, since we only discuss single-type users, we associate each category \( a \in [K] \) with its two unique parameters \( P_a := P_{a,1}, A_a := A_{a,1} \) and refer to them as scalars rather than vectors. In addition, we use the notation \( N_a \) for the random variable representing the number of iterations until a random user departs after being recommended by \( \pi^* \), the fixed-arm policy that recommends category \( a \) in each iteration.

To derive a regret bound for single-type users, we use two main lemmas: Lemma 4.1 which shows the optimal policy is fixed, and Lemma 4.2 which shows that returns of fixed-arm policies are sub-exponential and calculate their corresponding parameters. These lemmas allow us to use Algorithm 2 with a policy set \( \Pi \) that contains all the fixed-arm policies, and derive a \( \tilde{\tilde{O}}(\sqrt{T}) \) regret bound. For brevity, we relegate all the proofs to the Appendix.

To show that there exists a category \( a^* \in [K] \) for which \( \pi^{a^*} \) is optimal, we rely on the assumption that all the users have the same type (hence we drop the type subscripts \( t \)), and as a result the rewards of each category \( a \in [K] \) have an expectation that depends on a single parameter, namely \( \mathbb{E}[r(a)] = P_a \). Such a category \( a^* \) does not necessarily have the maximal click-probability nor the minimal departure-probability, but rather an optimal combination of the two (in a way, this is similar to the knapsack problem,
where we want to maximize the reward while having as little weight as possible. We formalize it in the following lemma.

**Lemma 4.1.** A policy $\pi^*$ is optimal if

$$a^* \in \arg \max_{a \in [K]} \frac{P_a}{\Lambda_a(1 - P_a)}.$$  

As a consequence of this lemma, the planning problem for single-type users is trivial—the solution is a fixed-arm policy $\pi^*$ given in the lemma. However, without access to the model parameters, identifying $\pi^*$ requires learning. We proceed with a simple observation regarding the random variable $N(a)$ for every category $a \in [K]$, where we want to maximize the reward while having as little weight as possible. 

**Lemma 4.2.** For every $a \in [K]$ and every $\Lambda_a > 0$, the random variable $N(a)$ follows a geometric distribution with success probability parameter $\Lambda_a [1 - P_a]$ in $(0, 1 - \epsilon)$.

Using Observation 4.2 and previously known results (stated as Lemma B.2 in the appendix), we show that $N(a)$ is sub-exponential for all $a \in [K]$. Notice that return realizations are always upper bounded by the trajectory length; this implies that returns are also sub-exponential. However, to use the regret bound of Algorithm 2, we need information regarding the parameters $(\tau_a^*, b_a)$ for every policy $\pi^*$. We provide this information in the following Lemma 4.3.

**Lemma 4.3.** For each category $a \in [K]$, the centred random variable $V^{\pi^*} - \mathbb{E}[V^{\pi^*}]$ is sub-exponential with parameters $(\tau_a^*, b_a)$, such that

$$\tau_a = b_a = \frac{8\epsilon}{\ln(1 - \Lambda_a(1 - P_a))}.$$  

**Proof sketch.** We rely on the equivalence between the subexponentiality of a random variable and the bounds on its moments (Property 2 in Definition 3.1). We bound the expectation of the return $V^{\pi^*}$, and use Minkowski’s and Jensen’s inequalities to show in Lemma B.2 that $\mathbb{E}[|V^{\pi^*} - \mathbb{E}[V^{\pi^*}]|^p]^{1/p}$ is upper bounded by $-4/\ln(1 - \Lambda_a(1 - P_a))$ for every $a \in [K]$ and $p \geq 1$. Finally, we apply a normalization trick and bound the Taylor series of $\mathbb{E}[\exp(\gamma(V^{\pi^*} - \mathbb{E}[V^{\pi^*}])]}$ to obtain the result. 

An immediate consequence of Lemma 4.3 is that the parameters $\tau_a^* = \frac{8\epsilon}{\ln(1 - \tau_a)}$ and $\eta = 1$ are valid upper bounds for $\tau_a$ and $b_a/\tau_a^2$ for each $a \in [K]$ (i.e., $\forall a \in [K] : \tau_a^* \geq \tau_a$ and $\eta \geq b_a^2/\tau_a^2$). We can now derive a regret bound using Algorithm 2 and Theorem 5.2.

**Theorem 4.4.** For single-type users ($M = 1$), running Algorithm 2 with $\Pi = \{\pi^* : a \in [K]\}$ and $\tilde{\tau} = \frac{8\epsilon}{\ln(1 - \tau_a^*)}, \eta = 1$ achieves an expected regret of at most

$$\mathbb{E}[R_T] = O(\sqrt{KT \log T + K \log T}).$$

**5 Two User Types and Two Categories**

In this section, we consider cases with two user types ($M = 2$), two categories ($K = 2$) and departure-probability $\Lambda_{\tau \tau} = 1$ for every category $a \in [K]$ and type $\tau \in [M]$. Even in this relatively simplified setting, where users leave after the first “no-click”, planning is essential. To see this, notice that the event of a user clicking on a certain category provides additional information about the user, which can be used to tailor better recommendations; hence, algorithms that do not take this into account may suffer a linear regret. In fact, this is not just a matter of the learning algorithm at hand, but rather a failure of all fixed-arm policies; there are instances where all fixed-arm policies yield high regret w.r.t. the baseline defined in Equation (1). Indeed, this is what the example in Section 2.1 showcases. Such an observation suggests that studying the optimal planning problem is vital.

In Section 5.1 we introduce the partially observable MDP formulation of departing bandits along with notion of belief-category walk. We use this notion to provide a closed-form formula for policies’ expected return, which we use extensively later on. Next, in Section 5.2 we characterize the optimal policy, and show that we can compute it in constant time relying on the closed-form formula. This is striking, as generally computing optimal POMDP policies is computationally intractable since, e.g., the space of policies grows exponentially with the horizon. Conceptually, we show that there exists an optimal policy that depends on a belief threshold: It recommends one category until the posterior belief of one type, which is monotonically increasing, crosses the threshold, and then it recommends the other category. Finally, in Section 5.3 we leverage all the previously obtained results to derive a small set of threshold policies of size $O(\ln T)$ with corresponding sub-exponential parameters. Due to Theorem 5.2 this result implies a $O(\sqrt{T})$ regret.

**5.1 Efficient Planning**

To recap, we aim to find the optimal policy when the click-probability matrix and the prior over user types are known. Namely, given an instance in the form of $(P, q)$, our goal is to efficiently find the optimal policy.

For planning purposes, the problem can be modeled by an episodic POMDP, $(S, [K], O, Tr, P, \Omega, q, O)$. A set of states, $S = [M] \cup \{\bot\}$ that comprises all types $[M]$, along with a designated absorbing state $\bot$ suggesting that the user departed (and the episode terminated). $[K]$ is the set of the actions (categories), $O = \{\text{stay}, \text{depart}\}$ is the set of possible observations. The transition and observation functions, $Tr : S \times [K] \rightarrow S$ and $\Omega : S \times [K] \rightarrow O$ (respectively) satisfy $Tr(\bot, i) = \Omega(\text{depart}, i) = 1 - P_{i,a}$ and $Tr(\text{stay}, i, a) = \Omega(\text{stay}, i, a) = P_{i,a}$ for every type $i \in [M]$ and action $a \in [K]$. Finally, $P$ is the expected reward matrix, and $q$ is the initial state distribution over the $M$ types.

When there are two user types and two categories, the click-probability matrix is given by Table 2 where we note that the prior on the types holds $q_i = 1 - q_{i+}$, thus can be represented by a single parameter $q_+$. 

**Remark 5.1.** Without loss of generality, we assume that $P_{1,2} \geq P_{2,1}, P_{1,2}, P_{2,1}$ since one could always permute the matrix to obtain such a structure.

Since the return and number of iterations for the same policy is independent of the user index, we drop the subscript $t$ in the rest of this subsection and use 


Table 2: Click probabilities for two user types and two categories.

|       | Type $x$ | Type $y$ |
|-------|----------|----------|
| Category 1 | $P_{1,x}$ | $P_{1,y}$ |
| Category 2 | $P_{2,x}$ | $P_{2,y}$ |
| Prior    | $q_x$    | $q_y = 1 - q_x$ |

As is well-known in the POMDP literature (Kaelbling, Littman, and Cassandra 1998), the optimal policy $\pi^*$ and its expected return are functions of belief states that represent the probability of the state at each time. In our setting, the states are the user types. We denote by $b_j$ the belief that the state is (type) $x$ at iteration $j$. Similarly, $1 - b_j$ is the belief that the state is (type) $y$ at iteration $j$. Needless to say, once the state $\bot$ is reached, the belief over the type states $[M]$ is irrelevant, as users do not come back. Nevertheless, we neglect this case as our analysis does not make use it.

We now describe how to compute the belief. At iteration $j = 1$, the belief state is set to be $b_1 = P(\text{state} = x) = q_x$. At iteration $j > 1$, upon receiving a positive reward $r_j = 1$, the belief is updated from $b_{j-1} \in [0,1]$ to

$$b_j(b_{j-1}, a, 1) = \frac{b_{j-1} \cdot P_{a,x}}{b_{j-1} \cdot P_{a,x} + P_{a,y}(1 - b_{j-1})},$$

(2)

where we note that in the event of no-click, the current user departs the system, i.e., we move to the absorbing state $\bot$. For any policy $\pi : [0,1] \rightarrow \{1,2\}$ that maps a belief to a category, its expected return satisfies the Bellman equation,

$$E[V^\pi(b)] = \left( b \cdot P_{\pi(b),x} + (1 - b) \cdot P_{\pi(b),y} \right) \cdot (1 + E[V^\pi(b', \pi(b), 1)]).$$

To better characterize the expected return, we introduce the following notion of belief-category walk.

**Definition 5.2 (Belief-category walk).** Let $\pi : [0,1] \rightarrow \{1,2\}$ be any policy. The sequence

$$b_1, a_1 = \pi(b_1), b_2, a_2 = \pi(b_2), \ldots$$

is called the belief-category walk. Namely, it is the induced walk of belief updates and categories chosen by $\pi$, given all the rewards are positive ($r_j = 1$ for every $j \in \mathbb{N}$).

Notice that every policy induces a single, well-defined and deterministic belief-category walk (recall that we assume departure-probabilities satisfy $A_{a,\bot} = 1$ for every $a \in [K], \tau \in [M]$). Moreover, given any policy $\pi$, the trajectory of every user recommended by $\pi$ is fully characterized by belief-category walk clipped at $b, a^{\pi}(t)$.

In what follows, we derive a closed-form expression for the expected return as a function of $b$, the categories chosen by the policy, and the click-probability matrix.

**Theorem 5.3.** For any policy $\pi$ and an initial belief $b \in [0,1]$, the expected return is given by

$$E[V^\pi(b)] = \sum_{i=1}^{\infty} b \cdot P_{1,x}^{m_{1,i}} \cdot P_{2,y}^{m_{2,i}} + (1 - b) \cdot P_{1,y}^{m_{1,i}} \cdot P_{2,y}^{m_{2,i}},$$

where $m_{1,i} := \sum_{j=1}^{i} a_j$ and $m_{2,i} := \sum_{j=1}^{i} \beta_j$ are calculated based on the belief-category walk $b_1, a_1, b_2, a_2, \ldots$ induced by $\pi$.

### 5.2 Characterizing the Optimal Policy

Using Theorem 5.3 we show that the planning problem can be solved in $O(1)$. To arrive at this conclusion, we perform a case analysis over the following three structures of the click-probability matrix $P$:

- **Dominant Row**, where $P_{1,y} \geq P_{2,y}$;
- **Dominant Column**, where $P_{1,x} \geq P_{2,x} > P_{1,y}$;
- **Dominant Diagonal**, where $P_{1,x} \geq P_{2,y} > P_{1,y}, P_{2,x}$.

Crucially, any matrix $P$ takes exactly one of the three structures. Further, since $P$ is known in the planning problem, identifying the structure at hand takes $O(1)$ time. Using this structure partition, we characterize the optimal policy.

**Dominant Row** We start by considering the simplest structure, in which the Category 1 is preferred by both types of users; Since $P_{1,y} \geq P_{2,y}$ and $P_{1,x} \geq P_{2,x}, P_{1,y}, P_{2,x}$ (Remark 5.1), there exists a dominant row, i.e., Category 1.

**Lemma 5.4.** For any instance such that $P$ has a dominant row $a$, the fixed policy $\pi^a$ is an optimal policy.

As expected, if Category 1 is dominant then the policy that always recommends Category 1 is optimal.

**Dominant Column** In the second structure we consider the case where there is no dominant row, and that the column of type $x$ is dominant, i.e., $P_{1,x} \geq P_{2,x} \geq P_{2,y} > P_{1,y}$. In such a case, which is also the one described in the example in Section 2.1, it is unclear what the optimal policy would be since none of the categories dominates the other.

Surprisingly, we show that the optimal policy can be of only one form: Recommend Category 2 for some time steps (possibly zero) and then always recommend Category 1. To identify when to switch from Category 2 to Category 1, one only needs to compare four expected returns.

**Theorem 5.5.** For any instance such that $P$ has a dominant column, one of the following four policies is optimal:

$$\pi^1, \pi^2, \pi^{2,[N^*]}, \pi^{2,[N^*]},$$

where $N^* = N^*(P, q)$ is a constant, and $\pi^{2,[N^*]} (\pi^{2,[N^*]}$) stands for recommending Category 2 until iteration $[N^*]$ ($[N^*]$) and then switching to Category 1.

The intuition behind the theorem is as follows. If the prior tends towards type $y$, we might start with recommending Category 2 (which users of type $y$ are more likely to click on). But after several iterations, and as long as the user stays, the posterior belief $b$ increases since $P_{2,x} > P_{2,y}$ (recall Equation 3). Consequently, since type $x$ becomes more probable, and since $P_{1,x} \geq P_{2,x}$, the optimal policy recommends the best category for this type, i.e., Category 1. For the exact expression of $N^*$, we refer the reader to Appendix 3.

Using Theorem 5.3 we can compute the expected return for each of the four policies in $O(1)$, showing that we can find the optimal policy when $P$ has a column in $O(1)$. 

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*References*

Kaelbling, Littman, and Cassandra (1998).
Dominant Diagonal. In the last structure, we consider the case where there is no dominant row (i.e., $P_{2,y} > P_{1,y}$) nor a dominant column (i.e., $P_{2,y} > P_{2,x}$). At first glance, this case is more complex than the previous two, since none of the categories and none of the types dominates the other one. However, we uncover that the optimal policy can be either always recommending Category 1 or always recommending Category 2. Theorem 5.6 summarizes this result.

**Theorem 5.6.** For any instance such that $P$ has a dominant diagonal, either $\pi^*$ or $\bar{\pi}$ is optimal.

With the full characterization of the optimal policy derived in this section (for all the three structures), we have shown that the optimal policy can be computed in $O(1)$.

5.3 Learning: UCB-based regret bound

In this section, we move from the planning task to the learning one. Building on the results of previous sections, we know that there must exist a threshold policy—a policy whose belief-category walk has a finite prefix of one category, and an infinite suffix with the other category—which is optimal. However, there can still be infinitely many such policies. To address this problem, we first show how to reduce the search space for approximately optimal policies with negligible additive factor to a set of $|\Pi| = O(ln(T))$ policies. Then, we derive the parameters $\tilde{\tau}$ and $\eta$ required for Algorithm 2. As an immediate consequence, we get a sublinear regret algorithm for this setting. We begin with defining threshold policies.

**Definition 5.7 (Threshold Policy).** A policy $\pi$ is called an $(a, h)$-threshold policy if there exists an number $h \in \mathbb{N} \cup \{0\}$ in $\pi$’s belief-category walk such that

- $\pi$ recommends category $a$ in iterations $j \leq h$, and
- $\pi$ recommends category $a'$ in iterations $j > h$,

for $a, a' \in \{1, 2\}$ and $a \neq a'$.

For instance, the policy $\pi^1$ that always recommends Category 1 is the $(2, 0)$-threshold policy, as it recommends Category 2 until the zero’th iteration (i.e., never recommends Category 2) and then Category 1 eternally. Furthermore, the policy $\pi^{2|[N^*]}$ introduced in Theorem 5.5 is the $(2, [N^*])$-threshold policy.

Next, recall that the chance of departure in every iteration is greater or equal to $\epsilon$, since we assume $\max_{a,\pi} P_{a,\pi} \leq 1 - \epsilon$. Consequently, the probability that a user will stay beyond $H$ iterations is exponentially decreasing with $H$. We could use high-probability arguments to claim that it suffices to focus on the first $H$ iterations, but without further insights this would yield $\Omega(2^H)$ candidates for the optimal policy. Instead, we exploit our insights about threshold policies.

Let $\Pi_H$ be the set of all $(a, h)$-threshold policies for $a \in \{1, 2\}$ and $h \in [H] \cup \{0\}$. Clearly, $|\Pi_H| = 2H + 2$. Lemma 5.8 shows that the return obtained by the best policy in $\Pi_H$ is not worse than that of the optimal policy $\pi^*$ by a negligible factor.

**Lemma 5.8.** For every $H \in \mathbb{N}$, it holds that

$$E\left[V^{\pi^*} - \max_{\pi \in \Pi_H} V^{\pi}\right] \leq \frac{1}{2O(H)}.$$
like what Algorithm 2 requires. More broadly, we could use this dynamic programming approach for more than two categories, namely for $K \geq 2$, but then the run-time becomes $O(H^K)$.

There are several interesting future directions. First, achieving low regret for the setup in Section 5 with $K \geq 2$. We suspect that this class of problems could enjoy a solution similar to ours, where candidates for optimal policies are mixing two categories solely. Second, achieving low regret for the setup in Section 5 with uncertain departure (i.e., $\mathbf{A} \neq 1$). Our approach fails in such a case since we cannot use belief-category walks; these are no longer deterministic. Consequently, the closed-form formula is much more complex and optimal planning becomes more intricate. These two challenges are left open for future work.

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A Extension: Planning Beyond Two User Types

In this section, we treat the planning task with two categories ($K = 2$) but potentially many types (i.e., $M \geq 2$). For convenience, we formalize the results in this section in terms of $M = 2$, but the results are readily extendable for the more general $2 \times M$ case. We derive an almost-optimal planning policy via dynamic programming, and then explain why it cannot be used for learning as we did in the previous section.

For reasons that will become apparent later on, we define by $V_H^\pi$ as the return of a policy $\pi$ until the $H$'s iteration. Using Theorem 5.3 we have that

$$E[V_H^\pi(b)] = \sum_{i=1}^{H} b \cdot \mathbf{P}_{1,x}^{m_{1,i}} \cdot \mathbf{P}_{2,x}^{m_{2,i}} + (1 - b) \mathbf{P}_{1,y}^{m_{1,i}} \cdot \mathbf{P}_{2,y}^{m_{2,i}},$$

where $m_{1,i} := |\{a_j = 1, j \leq i\}|$ and $m_{2,i} := |\{a_j = 2, j \leq i\}|$ are calculated based on the belief-category walk $b_1, a_1, b_2, a_2, \ldots$ induced by $\pi$. Further, let $\tilde{\pi}^*$ denote the policy maximizing $V_H^\pi$.

Notice that there is a bijection from $H$—iterations policies to $(m_{1,i}, m_{2,i})_{i=1}^{H}$; hence, we can find $\tilde{\pi}^*$ by finding the arg max of the expression on the right-hand-side of the above equation, in terms of $(m_{1,i}, m_{2,i})_{i=1}^{H}$. Formally, we want to solve the integer linear programming (ILP),

$$\begin{align*}
\text{maximize} & \quad \sum_{i=1}^{H} b \cdot \mathbf{P}_{1,x}^{m_{1,i}} \cdot \mathbf{P}_{2,x}^{m_{2,i}} + (1 - b) \mathbf{P}_{1,y}^{m_{1,i}} \cdot \mathbf{P}_{2,y}^{m_{2,i}} \\
\text{subject to} & \quad m_{a,i} = \sum_{l=1}^{i} z_{a,l} \text{ for } a \in \{1, 2\}, i \in [H], \\
& \quad z_{a,i} \in \{0, 1\} \text{ for } a \in \{1, 2\}, i \in [H], \\
& \quad z_{1,i} + z_{2,i} = 1 \text{ for } i \in [H].
\end{align*}$$

Despite that this problem involves integer programming, we can solve it using dynamic programming in $O(H^2)$ runtime. Notice that the optimization is over a subset of binary variables $(z_{1,i}, z_{2,i})_{i=1}^{H}$. Let $Z_H$ be the set of feasible solutions of the ILP, and similarly let $Z^H$ denote set of prefixes of length $h \leq H$ of $Z_H$.

For any $h \in [H]$ and $z \in Z^h$, define

$$D^h(z) \overset{\text{def}}{=} \sum_{i=1}^{h} b \cdot \mathbf{P}_{1,x}^{m_{1,i}} \cdot \mathbf{P}_{2,x}^{m_{2,i}} + (1 - b) \mathbf{P}_{1,y}^{m_{1,i}} \cdot \mathbf{P}_{2,y}^{m_{2,i}},$$

where $m_{a,i} = \sum_{l=1}^{i} z_{a,l}$ for $j \in \{1, 2\}, i \in [h]$ as in the ILP.

Consequently, solving the ILP is equivalent to maximizing $D^H$ over the domain $Z^H$.

Next, for any $h \in [H]$ and two integers $c_1, c_2$ such that $c_1 + c_2 = h$, define

$$\hat{D}^h(c_1, c_2) \overset{\text{def}}{=} \max_{z \in Z^h} D^h(z).$$

Under this construction, $\max_{c_1, c_2} \hat{D}^H(c_1, c_2)$ over $c_1, c_2$ such that $c_1 + c_2 = H$ is precisely the value of the ILP.

Reformulating Equation 3 for $h > 1$,

$$\hat{D}^h(c_1, c_2) = \max_{z_1 + z_2 = h} \{ \hat{D}^{h-1}(c_1 - z_1, c_2 - z_2) + \alpha(c_1, c_2) \},$$

where $\alpha(m_1, m_2) \overset{\text{def}}{=} b \cdot x_1^{m_1} \cdot x_2^{m_2} + (1 - b) y_1^{m_1} \cdot y_2^{m_2}$. For every $h$, there are only $h + 1$ possible values $\hat{D}^h$ can take: All the ways of dividing $h$ into non-negative integers $c_1$ and $c_2$; therefore, having computed $\hat{D}^{h-1}$ for all $h$ feasible inputs, we can compute $\hat{D}^h(c_1, c_2)$ in $O(h)$. Consequently, computing $\max_{c_1, c_2} \hat{D}^H(c_1, c_2)$, which is precisely the value of the ILP in 3, takes $O(H^2)$ run-time. Moreover, the policy $\tilde{\pi}^*$ can be found using backtracking. We remark that an argument similar to Lemma 5.8 implies that $E[V^* - V^{\tilde{\pi}^*}] \leq \frac{1}{2M+1}$, hence, $\tilde{\pi}^*$ is almost optimal.

To finalize this section, we remark that this approach could also work for $K > 2$ categories. Naively, for a finite horizon $H$, there are $K^H$ possible policies. The dynamic programming procedure explain above makes the search operate in run-time of $O(H^K)$. The run-time, exponential in the number of categories but polynomial in the horizon, is feasible when the number of categories is small.
B Experimental Evaluation

For general real-world datasets, we propose a scheme to construct semi-synthetic problem instances with many arms and many user types, using rating data sets with multiple ratings per user. We exemplify our scheme on the MovieLens Dataset [Harper and Konstan, 2015]. As a pre-processing step, we set movie genres to be the categories of interest, select a subset of categories \(|A|\) of size \(k\) (e.g., sci-fi, drama, and comedy), and select the number of user types, \(m\). Remove any user who has not provided a rating for at least one movie from each category \(a \in A\). When running the algorithm, randomly draw users from the data, and given a recommended category \(a\), suggest them a random movie which they have rated, and set their click probability to \(1 - r\), where \(r \in [0, 1]\) is their normalized rating of the suggested movie.

C UCB Policy for Sub-exponential Returns

An important tool for analyzing sub-exponential random variables is Bernstein’s Inequality, which is a concentration inequality for sub-exponential random variables (see, e.g., Jia, Shi, and Shen (2021)). Being a major component of the regret analysis for Algorithm 2, we state it here for convenience.

Lemma C.1. (Bernstein’s Inequality) Let a random variable \(X\) be sub-exponential with parameters \((\tau^2, b)\). Then for every \(v \geq 0\):
\[
\Pr[|X - \mathbb{E}[X]| \geq v] \leq \begin{cases} 
2 \exp\left(-\frac{v^2}{2\tau^2}\right) & v \leq \tau b \\
2 \exp\left(-\frac{v}{\tau}\right) & \text{else }
\end{cases}.
\]

D Single User Type: Proofs from Section 4

To simplify the proofs, we use the following notation: For a fixed-arm policy \(\pi\), we use \(V_j^\pi\) to denote its return from iteration \(j\) until the user departs. Namely,
\[
V_j^\pi = \sum_{i=j}^{N_\pi} \pi_a
\]
Throughout this section, we will use the following Observation.

Observation D.1. For every policy \(\pi\) and iteration \(j\),
\[
\mathbb{E}[V_j^\pi] = \pi_{\pi_j, 1} + \mathbb{E}[V_{j+1}^\pi] + (1 - \Lambda_{\pi_j})(1 - \pi_{\pi_j})\mathbb{E}[V_{j+1}^\pi] = \mathbb{E}[V_{j+1}^\pi](1 - \Lambda_{\pi_j}(1 - \pi_{\pi_j})) + \pi_{\pi_j}.
\]

Lemma 4.1. A policy \(\pi^{a^*}\) is optimal if
\[
a^* \in \arg\max_{a \in [K]} \frac{\pi_a}{\Lambda_a(1 - \pi_a)}.
\]

Proof. First, recall that every POMDP has an optimal Markovian policy which is deterministic (we refer the reader to Section 5.1 for full formulation of the problem as POMDP). Having independent rewards and a single state implies that there exists \(\mu^* \in \mathbb{N}\) such that \(\mathbb{E}[V_j^\pi] = \mu^*\) for every \(j \in \mathbb{N}\) (similarly to standard MAB problems, there exists a fixed-arm policy which is optimal).

Assume by contradiction that the optimal policy \(\pi^{a^*}\) holds
\[
a^* \notin \arg\max_{a \in [K]} \frac{\pi_a}{\Lambda_a(1 - \pi_a)}.
\]
Now, notice that
\[
\mathbb{E}[V^{a^*}] = \mathbb{E}[V_1^{a^*}] = \mathbb{E}[V_2^{a^*}](1 - \Lambda_{a^*}(1 - \pi_{a^*})) + \pi_{a^*}
\]
Solving the recurrence relation and summing the geometric series we get
\[
\mathbb{E}[V^{a^*}] = \pi_{a^*} \sum_{j=0}^{\infty} (1 - \Lambda_{a^*}(1 - \pi_{a^*}))^j = \frac{\pi_{a^*}}{\Lambda_{a^*}(1 - \pi_{a^*})}.
\]
Finally,
\[
a^* \notin \arg\max_{a \in [K]} \frac{\pi_a}{\Lambda_a(1 - \pi_a)},
\]
yields that any fixed-armed policy, \(\pi^{a'}\) such that
\[
a' \in \arg\max_{a \in [K]} \frac{\pi_a}{\Lambda_a(1 - \pi_a)}
\]
holds \(\mathbb{E}[V^{a^*}] > \mathbb{E}[V^{a^*}]\), a contradiction to the optimality of \(\pi^{a^*}\). □

Lemma D.2. For each $a \in [k]$, the centered random return $V^{\pi^a} - \mathbb{E}[V^{\pi^a}]$ is sub-exponential with parameter $C_2 = -4 / \ln(1 - \Lambda_a(1 - P_a))$.

In order to show that returns of fixed-arm policies are sub-exponential random variables, we first show that the number of iterations of users recommended by fixed-arm policies is also a sub-exponential. For this purpose, we state here a lemma that implies that every geometric r.v. is a sub-exponential r.v.. The proof of the next lemma appears, e.g., in [Hillar and Wibisono (2018) (Lemma 4.3)].

Lemma D.3. Let $X$ be a geometric random variable with parameter $r \in (0, 1)$, so that:

$$\Pr[X = x] = (1 - r)^{x-1} r, \quad x \in \mathbb{N}.$$  

Then $X$ satisfies Property (2) from Definition 3.1. Namely, $X$ is sub-exponential with parameter $C_2 = -2 / \ln(1 - r)$. Formally,

$$\forall p \geq 0 : (\mathbb{E}[X^p])^{1/p} \leq \frac{-2}{\ln(1 - r)}.$$  

The lemma above and Observation 4.2 allow us to deduce that the variables $N_a$ are sub-exponential in the first part of the following Corollary (the case in which $\Lambda_a = 0$ follows immediately from definition.). The second part of the lemma follows directly from the equivalences between Properties (2) and (1) in Definition 3.1.

Corollary D.4. For each $a \in [K]$, the number of iterations a user recommended by $\pi^a$ stays within the system, $N_a$, is sub-exponential with parameter $C_2^a = -2 / \ln(1 - \Lambda_a(1 - P_a))$. In addition, there exist constants $C_1^a > 0$ for every $a \in [K]$ such that

$$\forall a \in [K], \ v \geq 0 : \Pr[|N_a| > v] \leq \exp(1 - \frac{v}{C_1}).$$

The next Proposition D.5 is used for the proof of Lemma D.2.

Proposition D.5. For every $a \in [K]$, $|\mathbb{E}[V^{\pi^a}]| \leq \frac{-2}{\ln(1 - \Lambda_a(1 - P_a))}$.

Proof. First, notice that

$$(1 - \Lambda_a(1 - P_a)) \ln(1 - \Lambda_a(1 - P_a)) > (1 - \Lambda_a(1 - P_a)) - \frac{\Lambda_a(1 - P_a)}{1 - \Lambda_a(1 - P_a)} = -\Lambda_a(1 - P_a) > -2\Lambda_a(1 - P_a),$$

where the first inequality is due to $\frac{x}{1 + x} \leq \ln(1 + x)$ for every $x \geq -1$. Rearranging,

$$\frac{1 - \Lambda_a(1 - P_a)}{\Lambda_a(1 - P_a)} \leq \frac{-2}{\ln(1 - \Lambda_a(1 - P_a))}.$$  

(5)

For each user, the realization of $V^{\pi^a}$ is less or equal to the realization of $N_a - 1$ for the same user (as users provide negative feedback in their last iteration); hence,

$$|\mathbb{E}[V^{\pi^a}]| = \mathbb{E}[V^{\pi^a}] \leq \mathbb{E}[N_a] - 1 = \frac{1}{\Lambda_a(1 - P_a)} - 1 = \frac{1 - \Lambda_a(1 - P_a)}{\Lambda_a(1 - P_a)} \leq \frac{-2}{\ln(1 - \Lambda_a(1 - P_a))}.$$  

We proceed by showing that returns of fixed-arm policies satisfy Property (1) from Definition 3.1.

Lemma D.2. For each $a \in [k]$, the centered random return $V^{\pi^a} - \mathbb{E}[V^{\pi^a}]$ is sub-exponential with parameter $C_2 = -4 / \ln(1 - \Lambda_a(1 - P_a))$.

Proof. We use Property (1) from Definition 3.1 to derive that $V^{\pi^a}$ is also sub-exponential. This is true since the tails of $V^{\pi^a}$ satisfy that for all $v \geq 0$,

$$\Pr[|V^{\pi^a}| > v] \leq \Pr[|N_a| > v + 1] \leq \Pr[|N_a| > v] \leq \exp(1 - \frac{v}{C_1}),$$

where the first inequality follows since $|N_a| > v + 1$ is a necessary condition for $|V^{\pi^a}| > v$, and the last inequality follows from Corollary D.4. Along with Definition 3.1 we conclude that

$$|\mathbb{E}[V^{\pi^a}]|^{1/p} \leq \frac{-2}{\ln(1 - \Lambda_a(1 - P_a))}.$$  

(6)

Now, applying Minkowski’s inequality and then Jensen’s inequality (as $f(z) = z^p$, $g(z) = |z|$ are convex for every $p \geq 1$) we get

$$|\mathbb{E}[V^{\pi^a} - \mathbb{E}[V^{\pi^a}]|^p|^{1/p} \leq |\mathbb{E}[V^{\pi^a}]|^p|^{1/p} + |\mathbb{E}[\mathbb{E}[V^{\pi^a}]|^p|^{1/p} \leq |\mathbb{E}[V^{\pi^a}]|^p|^{1/p} + |\mathbb{E}[V^{\pi^a}]|. $$
Using Proposition D.5 and Inequality (6), we get
\[
\mathbb{E}[|V^{π^n}|^{1/p} + |E[V^{π^n}]|] \leq \frac{-2}{\ln(1 - A_a(1 - P_a))} + \frac{1}{A_a(1 - P_a)} - 1 \leq \frac{-4}{\ln(1 - A_a(1 - P_a))}
\]
Hence \(V^{π^n} - E[V^{π^n}]\) is sub-exponential with parameter \(C_2 = -4/\ln(1 - A_a(1 - P_a))\).

**Lemma 4.3.** For each category \(a \in [K]\), the centred random variable \(V^{π^n} - E[V^{π^n}]\) is sub-exponential with parameters \((γ_a^2, b_a)\), such that
\[
τ_a = b_a = -\frac{8e}{\ln(1 - A_a(1 - P_a))}.
\]

**Proof.** Throughout this proof, we will use the sub-exponential norm, \(||·||_{ψ_1}\), which is defined as
\[
||Z||_{ψ_1} = \sup_{p \geq 1} \frac{(\mathbb{E}[|Z|^p])^{1/p}}{p}.
\]
Let
\[
X = \frac{V^{π^n} - E[V^{π^n}]}{||V^{π^n} - E[V^{π^n}]||_{ψ_1}}, \quad y = γ · ||V^{π^n} - E[V^{π^n}]||_{ψ_1}.
\]
We have that
\[
||X||_{ψ_1} = \left|\frac{V^{π^n} - E[V^{π^n}]}{||V^{π^n} - E[V^{π^n}]||_{ψ_1}}\right| = 1.
\]
Let \(γ\) be such that \(|γ| < 1/b_a = -\frac{\ln(1 - A_a(1 - P_a))}{8e}\). From Lemma D.2, we conclude that
\[
|γ| = \left|\frac{y}{||V^{π^n} - E[V^{π^n}]||_{ψ_1}}\right| \leq -\frac{\ln(1 - A_a(1 - P_a))}{8e} = \frac{1}{2e} · \left|\frac{1}{||V^{π^n} - E[V^{π^n}]||_{ψ_1}}\right|;
\]
hence, \(|y| < \frac{1}{2e}\).
Summing the geometric series, we get
\[
\sum_{p=2}^{∞} (e|y|)^p = \frac{e^2y^2}{1 - e|y|} < 2e^2y^2 \tag{8}
\]
In addition, notice that \(yX = γ(V^{π^n} - E[V^{π^n}])\).
Next, from the Taylor series of \(\exp(·)\) we have
\[
\mathbb{E}[\exp(γ(V^{π^n} - E[V^{π^n}]))] = \mathbb{E}[\exp(yX)] = 1 + y\mathbb{E}[x] + \sum_{p=2}^{∞} \frac{y^p\mathbb{E}[X^p]}{p!}.
\]
Combining the fact that \(E[X] = 0\) and (7) to the above,
\[
1 + y\mathbb{E}[x] + \sum_{p=2}^{∞} \frac{y^p\mathbb{E}[X^p]}{p!} \leq 1 + \sum_{p=2}^{∞} \frac{y^p\mathbb{E}[X^p]}{p!}.
\]
By applying \(p! \geq (\frac{p}{2})^p\) and (3), we get
\[
1 + \sum_{p=2}^{∞} \frac{y^p\mathbb{E}[X^p]}{p!} \leq 1 + \sum_{p=2}^{∞} (e|y|)^p \leq 1 + 2e^2y^2 \leq \exp(2e^2y^2) = \exp(2e^2(γ · ||V^{π^n} - E[V^{π^n}]||_{ψ_1})^2),
\]
where the last inequality is due to \(1 + x \leq e^x\).
Note that \(||V^{π^n} - E[V^{π^n}]||_{ψ_1} \leq -\frac{4}{\ln(1 - A_a(1 - P_a))^2}\). Ultimately,
\[
\mathbb{E}[\exp(γ(V^{π^n} - E[V^{π^n}]))] \leq \exp \left(2e^2γ^2(-\frac{4}{\ln(1 - A_a(1 - P_a))^2})^2\right) = \exp \left(\frac{1}{2}γ^2(-\frac{8e}{\ln(1 - A_a(1 - P_a))^2})^2\right).
\]
This concludes the proof of the lemma.
E Two User Types and Two Categories: Proofs from Section [5]

E.1 Planning when \( K = 2 \)

Theorem 5.3. For every policy \( \pi \) and an initial belief \( b \in [0, 1] \), the expected return is given by

\[
\mathbb{E}[V^\pi(b)] = \sum_{i=1}^{\infty} b \cdot P_{1,x}^{m_{1,i}} \cdot P_{2,x}^{m_{2,i}} + (1 - b) P_{1,y}^{m_{1,i}} \cdot P_{2,y}^{m_{2,i}},
\]

where \( m_{1,i} := |\{a_j = 1, j \leq i\}| \) and \( m_{2,i} := |\{a_j = 2, j \leq i\}| \) are calculated based on the belief-category walk \( b_1, a_1, b_2, a_2, \ldots \) induced by \( \pi \).

Proof. Let \( \beta_t^\pi(b) := b \cdot P_{1,x}^{m_{1,i}} \cdot P_{2,x}^{m_{2,i}} + (1 - b) P_{1,y}^{m_{1,i}} \cdot P_{2,y}^{m_{2,i}} \). We will prove that for every policy \( \pi \) and every belief \( b \), we have that \( \mathbb{E}[V_H^\pi(b)] = \sum_{i=1}^{H} \beta_t^\pi(b) \) by a backward induction over \( H \).

For the base case, consider \( H = 1 \). We have that

\[
\mathbb{E}[V_1^\pi(b_1)] = b_1 \cdot P_{a_1,x} + (1 - b_1) P_{a_1,y} = b_1 \cdot P_{1,x}^{m_{1,1}} \cdot P_{2,x}^{m_{2,1}} + (1 - b_1) P_{1,y}^{m_{1,1}} \cdot P_{2,y}^{m_{2,1}} = \beta_1^\pi(b)
\]
as \( m_{a,1} = 1 \).

For the inductive step, assume that \( \mathbb{E}[V_{H-1}^\pi(b)] = \sum_{i=1}^{H-1} \beta_t^\pi(b) \) for every \( b \in [0, 1] \). We need to show that \( \mathbb{E}[V_H^\pi(b)] = \sum_{i=1}^{H} \beta_t^\pi(b) \) for every \( b \in [0, 1] \).

Indeed,

\[
\mathbb{E}[V_H^\pi(b_1)] = \beta_1^\pi(b_1)(1 + \mathbb{E}[V_{H-1}^\pi(b_1)])
\]

\[
= \beta_1^\pi(b_1)(1 + \mathbb{E}[V_{H-1}^\pi(b_2)])
\]

\[
= \beta_1^\pi(b_1)(1 + \sum_{i=2}^{H-1} \beta_t^\pi(b_2))
\]

\[
= \sum_{i=1}^{H} \beta_t^\pi(b_1),
\]

where the second to last equality is due to the induction hypothesis and the assumption that \( \pi \) is a deterministic stationary policy. The proof completes by realizing that \( \mathbb{E}[V^\pi(b)] = \lim_{H \to \infty} \mathbb{E}[V_H^\pi(b)] = \lim_{H \to \infty} \sum_{i=1}^{H} \beta_t^\pi(b) = \sum_{i=1}^{\infty} \beta_t^\pi(b) \), since the sum is finite and has positive summands.

E.2 Dominant Row (DR)

Lemma 5.4. For any instance such that \( P \) has a dominant row \( a \), the fixed policy \( \pi^a \) is an optimal policy.

Proof. We will show that for every iteration \( j \), no matter what were the previous topic recommendations were, selecting topic 1 rather than topic 2 can only increase the value.

Let \( \pi \) be a stationary policy such that \( \pi(b_j) = 2 \). Changing it into a policy \( \pi^j \) that is equivalent to \( \pi \) for all iterations but iteration \( j + 1 \) in which it recommends topic 1 can only improve the value.

Since \( P_{1,x} > 0, P_{2,x} \geq 0, P_{1,y} \geq 0 \) and \( b, 1 - b \geq 0 \) and this structure satisfies \( P_{2,y} - P_{1,y} \leq 0 \), we get that for every \( \tilde{m}_{1,j}, \tilde{m}_{2,j}, n_{1,j}, n_{2,j} \in \mathbb{N} \) and for every \( b \),

\[
b \cdot P_{1,x}^{\tilde{m}_{1,j} + n_{1,j}} \cdot P_{2,x}^{\tilde{m}_{2,j} + n_{2,j}} (P_{1,x} - P_{2,x}) \geq (1 - b) P_{1,y}^{\tilde{m}_{1,j} + n_{1,j}} \cdot P_{2,y}^{\tilde{m}_{2,j} + n_{2,j}} (P_{2,y} - P_{1,y});
\]

thus,

\[
b \cdot P_{1,x}^{\tilde{m}_{1,j} + 1 + n_{1,j}} \cdot P_{2,x}^{\tilde{m}_{2,j} + 1 + n_{2,j}} + (1 - b) P_{1,y}^{\tilde{m}_{1,j} + 1 + n_{1,j}} \cdot P_{2,y}^{\tilde{m}_{2,j} + 1 + n_{2,j}} \geq \]

\[
b \cdot P_{1,x}^{\tilde{m}_{1,j} + n_{1,j}} \cdot P_{2,x}^{\tilde{m}_{2,j} + n_{2,j}} + (1 - b) P_{1,y}^{\tilde{m}_{1,j} + n_{1,j}} \cdot P_{2,y}^{\tilde{m}_{2,j} + n_{2,j}}.
\]

Hence for every time step \( j + 1 \), choosing topic 1 instead of topic 2 leads to increased value of each of the summation element \( b \cdot P_{1,x}^{m_{1,i}} \cdot P_{2,x} + (1 - b) P_{1,y}^{m_{1,i}} \cdot P_{2,y}^{m_{2,i}} \) such that \( m_{1,i} = \tilde{m}_{1,j} + n_{1,j} \geq \tilde{m}_{1,j} \) and \( m_{2,i} = \tilde{m}_{2,j} + n_{2,j} \geq \tilde{m}_{2,j} \). We deduce that

\[
\mathbb{E}[V^{\pi^j}(b)] \geq \mathbb{E}[V^\pi(b)].
\]
E.3 Dominant Column (DC)

Before proving the main theorem (Theorem 5.5), we prove two auxiliary lemmas.

Lemma E.1. For \( P \) with a DC structure, if a policy \( \pi \) is optimal then it recommends topic 1 for all iteration \( j' \geq j + 1 \) such that

\[
\sum_{i=j+1}^{\infty} b \cdot P_{1,x}^{m_{1,i}^{1}} \cdot P_{2,x}^{m_{2,i}^{1}} > \sum_{i=j+1}^{\infty} \frac{1 - b}{b} \cdot P_{2,y} - \frac{P_{1,y}}{P_{1,x} - P_{2,x}} \cdot \frac{P_{2,y} m_{1,i}^{1} + P_{2,y} m_{2,i}^{1}}{P_{2,y} 1 + y + P_{2,y}} .
\]  

(9)

Proof. First, assume by contradiction that there exists an optimal policy \( \pi \) that recommends topic 2 in iteration \( j + 1 \) such that (9) holds.

Let \( \pi^j \) be the policy that is equivalent to \( \pi \) but recommend topic 1 instead of topic 2 in iteration \( j + 1 \). Since \( \pi \) and \( \pi^j \) recommends the same topic until iteration \( j \), along with the optimality of \( \pi \), we have

\[
\mathbb{E}[V_{\pi^j}(b)] - \mathbb{E}[V_{\pi}(b)] = \mathbb{E}[V_{\pi^j}(b)] - \mathbb{E}[V_{\pi}(b)] \leq 0.
\]

Expanding the above equation,

\[
\sum_{i=j+1}^{\infty} b \cdot P_{1,x}^{m_{1,i}^{1}} \cdot P_{2,x}^{m_{2,i}^{1} - 1} + (1 - b)P_{1,y}^{m_{1,i}^{1} - 1} \cdot \sum_{i=j+1}^{\infty} b \cdot P_{1,x}^{m_{1,i}^{1}} \cdot P_{2,x}^{m_{2,i}^{1}} + (1 - b)P_{1,y} \cdot P_{2,y} \leq 0
\]

\[
\sum_{i=j+1}^{\infty} b \cdot P_{1,x}^{m_{1,i}^{1}} \cdot P_{2,x}^{m_{2,i}^{1}} \leq \sum_{i=j+1}^{\infty} P_{2,x} - \sum_{i=j+1}^{\infty} (\frac{P_{2,y} - P_{1,y}}{P_{2,y}}) \sum_{i=j+1}^{\infty} P_{1,x} \cdot P_{2,x} - \sum_{i=j+1}^{\infty} P_{1,y} \cdot P_{2,y}
\]

which is a contradiction to (9).

For the second part of the lemma, assume that condition (9) holds for some iteration \( j + 1 \in \mathbb{N} \) and some optimal policy \( \pi \), hence, \( \pi(b, m_{1,j}^{1}, m_{2,j}^{1}) = 1 \) and we have \( m_{1,j+1}^{1} = m_{1,j}^{1} + 1 \) and \( m_{2,j+1}^{1} = m_{2,j}^{1} \). Exploiting this fact, we have that

\[
\sum_{i=j+2}^{\infty} P_{1,x}^{m_{1,i}^{1}} P_{2,x}^{m_{2,i}^{1}} = \sum_{i=j+1}^{\infty} P_{1,x}^{m_{1,i}^{1} + 1} P_{2,x}^{m_{2,i}^{1}} = \sum_{i=j+1}^{\infty} P_{1,x}^{m_{1,i}^{1}} P_{2,x}^{m_{2,i}^{1}} > (9)
\]

implying

\[
(P_{1,x} \geq P_{1,y}) P_{1,y} \sum_{i=j+1}^{\infty} \frac{1 - b}{b} \cdot P_{2,y} - P_{1,y} \cdot P_{2,x} P_{1,y}^{m_{1,i}^{1}} P_{2,y}^{m_{2,i}^{1}}
\]

\[
\geq (P_{1,x} \geq P_{1,y}) P_{1,y} \sum_{i=j+1}^{\infty} \frac{1 - b}{b} \cdot P_{2,y} - P_{1,y} \cdot P_{2,x} P_{1,y}^{m_{1,i}^{1} + 1} P_{2,y}^{m_{2,i}^{1}}
\]

\[
= \sum_{i=j+1}^{\infty} \frac{1 - b}{b} \cdot P_{2,y} - P_{1,y} \cdot P_{2,x} P_{1,y}^{m_{1,i}^{1}} P_{2,y}^{m_{2,i}^{1}}
\]

\[
= \sum_{i=j+2}^{\infty} \frac{1 - b}{b} \cdot P_{2,y} - P_{1,y} \cdot P_{2,x} P_{1,y}^{m_{1,i}^{1}} P_{2,y}^{m_{2,i}^{1}}
\]

An immediate consequence of Lemma E.1 is the following corollary.

Corollary E.2. For any DC-structured \( P \) and every belief \( b \in [0, 1] \), the optimal policy \( \pi \) first recommends topic 2 for at most

\[
\arg\min_{N} \sum_{i=1}^{N} P_{2,x}^{m_{2,i}^{1}} > \frac{1 - b}{b} \cdot P_{2,y} - P_{1,y} \cdot P_{2,x} \sum_{i=1}^{N} P_{2,y}^{m_{2,i}^{1}}
\]

times, and then recommends topic 1 permanently. In addition, \( N \in \mathbb{N} \) since \( P_{2,x} > P_{2,y} \).
Theorem 5.5. For any instance such that \( \textbf{P} \) has a dominant column, one of the following four policies is optimal:
\[
\pi^1, \pi^2, \pi^1[\text{DD}], \pi^2[\text{DD}],
\]
where \( N^* = N^*(\textbf{P}, \textbf{q}) \) is a constant, and \( \pi^1[\text{DD}], \pi^2[\text{DD}] \) stand for recommending Category 2 until iteration \( \lfloor N^* \rfloor \) and then switching to Category 1.

Proof. Due to Theorem 5.3 and Corollary E.2 we can write the expected value of a policy as a function of \( \pi \)

\[
\begin{align*}
\mathbb{E}[V^{\pi N}(b)] &= \sum_{i=1}^{b} \cdot \textbf{P}_{1,x}^{m_1,i} \cdot \textbf{P}_{2,x}^{m_2,i} + (1-b)\textbf{P}_{1,y}^{m_1,i} \cdot \textbf{P}_{2,y}^{m_2,i} \\
&= \sum_{b=1}^{N} \cdot \textbf{P}_{2,x}^{i} + (1-b)\textbf{P}_{2,y}^{i} + \sum_{i=1}^{\infty} \cdot \textbf{P}_{1,x}^{i-N} \cdot \textbf{P}_{1,y}^{i-N} \\
&= b \cdot \textbf{P}_{2,x}(\textbf{P}_{2,x}^{N} - 1) + (1-b) \cdot \textbf{P}_{2,y}(\textbf{P}_{2,y}^{N} - 1) + b \cdot \textbf{P}_{2,x}^{N} \cdot (1-b)\textbf{P}_{2,y}^{N} \\
&= b \cdot \textbf{P}_{2,x}(\textbf{P}_{2,x}^{N} - 1) + (1-b) \cdot \textbf{P}_{2,y}(\textbf{P}_{2,y}^{N} - 1) + b \cdot \textbf{P}_{2,x}^{N} \cdot (1-b)\textbf{P}_{2,y}^{N}.
\end{align*}
\]

Equation (10) could be cast as \( c_1 \cdot \textbf{P}_{2,x}^{N} + c_2\textbf{P}_{2,y}^{N} + c_3(\textbf{P}_{2,x}, \textbf{P}_{2,y}) \) for positive \( c_1 \), negative \( c_2 \) and positive \( c_3 \). Let \( f : \mathbb{R} \rightarrow \mathbb{R} \) be the continuous function such that \( f(N) = c_1 \cdot \textbf{P}_{2,x}^{N} + c_2\textbf{P}_{2,y}^{N} + c_3(\textbf{P}_{2,x}, \textbf{P}_{2,y}) \).

We take the derivative w.r.t. \( N \) to find the saddle point of \( f \):
\[
\frac{d}{dN} f = c_1 \cdot \ln \textbf{P}_{2,x} \cdot \textbf{P}_{2,x}^{N} + c_2 \ln \textbf{P}_{2,y} \cdot \textbf{P}_{2,y}^{N} = 0,
\]
which suggests the saddle point of \( f \) is
\[
N^* = \frac{\ln \left( -\frac{c_2 \ln \textbf{P}_{2,x}}{c_1 \ln \textbf{P}_{2,x}} \right)}{\ln (\frac{\textbf{P}_{2,x}}{\textbf{P}_{2,y}})}.
\]

Next, set \( N^{*} \overset{\text{df}}{=} \max \{0, N^* \} \). Since \( f \) has a single saddle point and for every \( n \in \mathbb{N} \) it holds that \( f(N) = \mathbb{E}[V^{\pi N}(b)] \),

to determine the optimal policy, one only needs to compare the value \( \mathbb{E}[V^{\pi N}(b)] \) at the boundary points \( (N = 0, N = \infty) \) and at the closest integers to the saddle point \( (N = [N^*], N = [N^{*}]) \).

---

E.4 Dominant Diagonal (SD)

Theorem 5.6. For any instance such that \( \textbf{P} \) has a dominant diagonal, either \( \pi^1 \) or \( \pi^2 \) is optimal.

Proof. We prove the following claim by a backward induction over the number of iterations remaining: For every \( k = H - 1, \ldots, 1 \) it holds that for every policy \( \pi \) and belief \( b \),
\[
\mathbb{E}[V^{\pi N}(b)] \leq \max \{\mathbb{E}[V^{\pi_1 N}(b)], \mathbb{E}[V^{\pi_2 N}(b)]\}.
\]

First, we notice that when \( k = H - 1 \), the only possible policies are \( \pi^1 \) and \( \pi^2 \). For \( k = H - 2 \), we prove the statement by contradiction. There are only two ways to select topics when \( k = H - 2 \):
\[
\pi'_1 = (\pi_{1:H-2}, \frac{1}{H-1}, \frac{2}{H}) \quad \text{and} \quad \pi''_1 = (\pi_{1:H-2}, \frac{2}{H-1}, \frac{1}{H}).
\]

Let \( m_1 \) and \( m_2 \) denote the number of times \( \pi \) has played topic 1 and 2 till time \( H - 2 \), inclusive. Assume that the policy \( \pi' \) is optimal. In particular, it holds that \( \mathbb{E}[V^{\pi_1}] \leq \mathbb{E}[V^{\pi_2}] \) and \( \mathbb{E}[V^{\pi_2}] \leq \mathbb{E}[V^{\pi_1}] \). We get
\[
\begin{align*}
\text{and}
\end{align*}
\]

\[
\begin{align*}
\mathbb{E}[V^{\pi_1 N}(b)] &\leq \mathbb{E}[V^{\pi_2 N}(b)] \\
&\leq \mathbb{E}[V^{\pi_1 N}(b)],
\end{align*}
\]

(11) and
(12)
As both sides of (11) and (12) are positive, the multiplication of their left hand sides is smaller than the multiplication of their right hand sides, i.e.,

\[
bP_{1,x}^mP_{2,x}b(P_{1,x} - P_{2,x})P_{1,y}^mP_{2,y}^m(1 - b)(P_{2,y} - P_{1,y})(1 + P_{2,y}) \\
\leq P_{1,y}^mP_{2,y}^m(1 - b)P_{1,y}(P_{2,y} - P_{1,y})bP_{1,x}^mP_{2,x}^m(P_{1,x} - P_{2,x})(1 + P_{2,x})
\]

Dividing both sides by \(bP_{1,x}^mP_{2,x}^m(P_{1,x} - P_{2,x})P_{1,y}^mP_{2,y}^m(1 - b)(P_{2,y} - P_{1,y}) > 0\), we obtain

\[P_{1,x}(1 + P_{2,y}) \leq P_{1,y}(1 + P_{2,x}),\]

which is a contradiction as \(P_{1,x} > P_{1,y}\) and \(1 + P_{2,y} > 1 + P_{2,x}\).

Now, assume that the policy \(\pi''\) is optimal. In particular, it holds that \(E[V_{\pi''}^t] \leq E[V_{\pi'}^t]\) and \(E[V_{\pi''}^t] \leq E[V_{\pi'}^t]\). We get

\[P_{1,x}^mP_{2,x}^m(P_{1,x} - P_{2,x})(1 + P_{1,x}) \leq P_{1,y}^mP_{2,y}^m(1 - b)(1 + P_{1,y})(P_{2,y} - P_{1,y}),\] (13)

and

\[P_{1,y}^mP_{2,y}^m(1 - b)P_{2,y}(P_{2,y} - P_{1,y}) \leq P_{1,x}^mP_{2,x}^mP_{2,x}(P_{1,x} - P_{2,x}).\] (14)

As both sides of (13) and (14) are positive, the multiplication of their left hand sides is smaller than the multiplication of their right hand sides,

\[P_{1,x}^mP_{2,x}^m(P_{1,x} - P_{2,x})(1 + P_{1,x})P_{1,y}^mP_{2,y}^m(1 - b)P_{2,y}(P_{2,y} - P_{1,y}) \\
\leq P_{1,y}^mP_{2,y}^m(1 - b)(1 + P_{1,y})(P_{2,y} - P_{1,y})P_{1,x}^mP_{2,x}^mP_{2,x}(P_{1,x} - P_{2,x}).\]

Dividing both sides by \(P_{1,x}^mP_{2,x}^m(P_{1,x} - P_{2,x})P_{1,y}^mP_{2,y}^m(1 - b)(P_{2,y} - P_{1,y}) > 0\), we obtain

\[P_{2,y}(1 + P_{1,x}) \leq P_{2,x}(1 + P_{1,y}),\]

which is again, a contradiction as \(P_{2,x} < P_{2,y}\) and \(1 + P_{1,y} < 1 + P_{1,x}\).

For \(H \geq 3\), we prove the statement by contradiction. Suppose not, i.e., the optimal policy \(\pi\) switch recommended topic at least once. Let \(t\) denote the time step where \(\pi\) switch for the last time. We first consider the case where \(\pi\) has switched from topic 2 to topic 1 at time \(t\). More specifically, we have

\[\pi_{1:H} = (\pi_{1:t-2}, \frac{2}{\pi_{t-1}}, \frac{1}{\pi_1}, \frac{1}{\pi_{t+1:H-1}}, \frac{1}{\pi_H}).\]

Consider another policy \(\tilde{\pi}\) (that behaves the same as \(\pi\) except at time step \(t - 1\)) defined as

\[\tilde{\pi}_{1:H} = (\pi_{1:t-2}, \frac{2}{\pi_{t-1}}, \frac{2}{\pi_1}, \frac{1}{\pi_{t+1:H-1}}, \frac{1}{\pi_H}).\]

Let \(m_1\) and \(m_2\) denote the number of times \(\pi\) has recommended topic 1 and 2 till (and include) time \(t - 1\). Since \(\pi\) is optimal, we have the difference between the value of \(\pi\) and \(\tilde{\pi}\) to be non-negative, i.e.,

\[E[V_{\pi'}^t] - E[V_{\tilde{\pi}}^t] = \sum_{i=1}^{H-t+1} bP_{1,x}^{m_1+i-1}P_{2,x}^{m_2+1}(P_{1,x} - P_{2,x}) + (1 - b)P_{1,y}^{m_1+i-1}P_{2,y}^{m_2+1}(P_{1,y} - P_{2,y}) \geq 0,\] (15)

where the difference is induced by the discrepancy of the two policies from time step \(t\) to \(H\). Consider another policy \(\pi'\) (that behaves the same as \(\pi\) except at time step \(H\)) defined as

\[\pi'_{1:H} = (\pi_{1:t-2}, \frac{2}{\pi_{t-1}}, \frac{1}{\pi_1}, \frac{1}{\pi_{t+1:H-1}}, \frac{2}{\pi_H}).\]

Since \(\pi\) is optimal, we have the difference between the value of \(\pi\) and \(\pi'\) to be non-negative, i.e.,

\[E[V_{\pi'}^t] > E[V_{\pi'}^t] \Rightarrow bP_{1,x}^{m_1+H-t}P_{2,x}^{m_2}(P_{1,x} - P_{2,x}) > (1 - b)P_{1,y}^{m_1+H-t}P_{2,y}^{m_2}(P_{2,y} - P_{1,y}),\]

where the difference is induced by the discrepancy of the two policies from time step \(H\). Multiplying both sides by \(P_{1,y} > 0\), we get

\[P_{1,y}bP_{1,x}^{m_1+H-t}P_{2,x}^{m_2}(P_{1,x} - P_{2,x}) > (1 - b)P_{1,y}^{m_1+H-t+1}P_{2,y}^{m_2}(P_{2,y} - P_{1,y}).\]

Using \(P_{1,x}^{m_1+H-t}P_{2,x}^{m_2}(P_{1,x} - P_{2,x}) > 0\), and \(P_{1,y}bP_{1,x}^{m_1+H-t}P_{2,x}^{m_2}(P_{1,x} - P_{2,x}) > 0,\)

\[bP_{1,x}^{m_1+H-t+1}P_{2,x}^{m_2}(P_{1,x} - P_{2,x}) > (1 - b)P_{1,y}^{m_1+H-t+1}P_{2,y}^{m_2}(P_{2,y} - P_{1,y});\]
where the difference follows from the discrepancy between the two policies from time step \( t \) to \( \tau \):

\[
bP_{1,x}^{m_1 + H - t + 1}P_{2,x}^{m_2}P_{1,y} - (1 - b)P_{1,y}^{m_1 + H - t + 1}P_{2,y}^{m_2}P_{1,y} - P_{2,y} \geq 0. \tag{16}
\]

Next, we construct a new policy \( \pi_{\text{new}} \) that outperforms \( \pi \). We let \( \pi_{\text{new}}^{1:H} \) to be the policy defined as below

\[
\pi_{\text{new}}^{1:H} = (\pi_{1:t-1}, 1, 1, \ldots, 1, \pi_{t+1:H}).
\]

The value difference between \( \pi_{\text{new}}^{1:H} \) and \( \pi \) (caused by the discrepancy of the two policies from time \( t \) to \( H \)) is

\[
E[V_{\pi_{\text{new}}}^{1:H}] - E[V_{\pi}^{1:H}] = \sum_{i=1}^{H-t+1} bP_{1,x}^{m_1 + i - 1}P_{2,x}^{m_2}P_{1,y} - P_{2,y} + (1 - b)P_{1,y}^{m_1 + i - 1}P_{2,y}^{m_2}P_{1,y} - P_{2,y}
\]

\[
+ bP_{1,x}^{m_1 + H - t + 1}P_{2,x}^{m_2}P_{1,y} - P_{2,y} + (1 - b)P_{1,y}^{m_1 + H - t + 1}P_{2,y}^{m_2}P_{1,y} - P_{2,y}
\]

\[
> \sum_{i=1}^{H-t+1} bP_{1,x}^{m_1 + i - 1}P_{2,x}^{m_2 + 1}(P_{1,y} - P_{2,y}) + (1 - b)P_{1,y}^{m_1 + i - 1}P_{2,y}^{m_2 + 1}(P_{1,y} - P_{2,y})
\]

\[
\geq 0,
\]

where the first inequality is true because \( P_{2,x} < P_{2,y}, P_{1,x} - P_{2,x} > 0 \) and \( P_{1,y} - P_{2,y} < 0 \), therefore for every \( 1 \leq i \leq H - t + 1 \)

\[
bP_{1,x}^{m_1 + i - 1}P_{2,x}^{m_2}(P_{1,y} - P_{2,y})(1 - P_{2,y}) > 0 > (1 - b)P_{1,y}^{m_1 + i - 1}P_{2,y}^{m_2}(P_{2,y} - P_{1,y})(P_{2,y} - 1)
\]

along with (16). The second inequality follows from (15). Thus, we have successfully find another policy \( \pi_{\text{new}}^{1:H} \) that differs from \( \pi \) and achieves a higher value, which is a contradiction.

Next, we consider the case where \( \pi \) has switched from topic 1 to topic 2 at time \( t \), i.e.,

\[
\pi_{1:H} = (\pi_{1:t-1}, 1, 1, \ldots, 1, \pi_{t+1:H}).
\]

Consider another policy \( \tilde{\pi} \) (that behaves the same as \( \pi \) except at time step \( t \)) defined as

\[
\tilde{\pi}_{1:H} = (\pi_{1:t-2}, 1, 1, \ldots, 1, \pi_{t+1:H}).
\]

Since \( \pi \) is optimal, we have the difference between the value of \( \pi \) and \( \tilde{\pi} \) to be non-negative, i.e.,

\[
E[V_{\pi}^{1:H}] - E[V_{\tilde{\pi}}^{1:H}] = \sum_{i=1}^{H-t+1} bP_{1,x}^{m_1 + i - 1}P_{2,x}^{m_2 + i - 1}(P_{2,x} - P_{1,x}) + (1 - b)P_{1,y}^{m_1 + i - 1}P_{2,y}^{m_2 + i - 1}(P_{2,y} - P_{1,y}) \geq 0,
\]

where the difference follows from the discrepancy between the two policies from time step \( t \) to \( H \).

Consider another policy \( \pi' \) (that behaves the same as \( \pi \) except at time step \( H \)) defined as

\[
\pi'_{1:H} = (\pi_{1:t-2}, 1, 1, \ldots, 1, \pi_{t+1:H}).
\]

Since \( \pi \) is optimal, we have the difference between the value of \( \pi \) and \( \pi' \) to be non-negative, i.e.,

\[
E[V_{\pi}^{1:H}] > E[V_{\pi'}^{1:H}] \Rightarrow (1 - b)P_{1,y}^{m_1}P_{2,y}^{m_2 + H - t}(P_{2,y} - P_{1,y}) \geq bP_{1,x}^{m_1}P_{2,x}^{m_2 + H - t}(P_{1,x} - P_{2,x}),
\]

where the difference is induced by the discrepancy of the two policies from time step \( H \). Multiplying both sides by \( P_{2,x} > 0 \),

\[
P_{2,x}(1 - b)P_{1,y}^{m_1}P_{2,y}^{m_2 + H - t}(P_{2,y} - P_{1,y}) \geq bP_{1,x}^{m_1}P_{2,x}^{m_2 + H - t + 1}(P_{1,x} - P_{2,x}).
\]

Using \( P_{2,x}(1 - b)P_{1,y}^{m_1}P_{2,y}^{m_2 + H - t}(P_{2,y} - P_{1,y}) > 0 \) and \( P_{2,x} \geq 1 \), we get

\[
(1 - b)P_{1,y}^{m_1}P_{2,y}^{m_2 + H - t + 1}(P_{2,y} - P_{1,y}) \geq bP_{1,x}^{m_1}P_{2,x}^{m_2 + H - t + 1}(P_{1,x} - P_{2,x});
\]

hence,

\[
bP_{1,x}^{m_1}P_{2,x}^{m_2 + H - t + 1}(P_{2,x} - P_{1,x}) + (1 - b)P_{1,y}^{m_1}P_{2,y}^{m_2 + H - t + 1}(P_{2,y} - P_{1,y}) \geq 0.
\]
Again, we will construct a new policy $\pi_{\text{new}}$ that outperforms $\pi$. We let $\pi_{\text{new}}$ to be the policy defined as below

$$
\pi_{\text{new}}(1) = \begin{cases} 
\pi_1, & \text{if } \pi_1 \geq \frac{2}{\pi_{t+1}} \pi_{t+1} + \ldots + \frac{2}{\pi_t}, \\
\pi_{t+1:H-1}, & \text{otherwise}.
\end{cases}
$$

Now, the value difference between $\pi_{\text{new}}$ and $\pi$ (caused by the discrepancy of the two policies from time $t-1$ to $H$) is

$$
\mathbb{E}[V_{\pi_{\text{new}}}^H] - \mathbb{E}[V_{\pi}^H] = \sum_{i=1}^{H-t+1} \left( bP_{1,x}^mP_{2,x}^{m+i-1}(P_{2,x} - P_{1,x}) + (1-b)P_{1,y}^mP_{2,y}^{m+i-1}(P_{2,y} - P_{1,y}) \right)
$$

$$
\geq \sum_{i=1}^{H-t+1} bP_{1,x}^mP_{2,x}^{m+i-1}(P_{2,x} - P_{1,x}) + (1-b)P_{1,y}^mP_{2,y}^{m+i-1}(P_{2,y} - P_{1,y})
$$

where the first inequality is true because $P_{1,y} < P_{1,x}$, $P_{2,x} - P_{1,x} < 0$ and $P_{2,y} - P_{1,y} > 0$ and (18), and the second from (17). Similarly, we have successfully find another policy $\pi_{\text{new}}^{\tau}$ that differs from $\pi$ and achieves a higher value, which is a contradiction.

We have covered all cases, so the inductive argument holds. This concludes the proof of the theorem.

\section{E.5 UCB-based regret bound}

\textbf{Lemma 5.9.} Let $\hat{\tau} = \frac{8e}{\ln\left(\frac{1}{1-\epsilon}\right)}$ and $\eta = 1$. For every threshold policy $\pi \in \Pi_H$, the centred random variable $V^\pi - \mathbb{E}[V^\pi]$ is $(\tau^2, b)$-sub-exponential with $(\tau^2, b)$ satisfying $\hat{\tau} \geq \tau$ and $\eta \geq b^2/\tau^2$.

\textbf{Proof.} Let $\gamma$ be such that

$$
|\gamma| < -\frac{\ln(1-\epsilon)}{8e} \leq \min_{a\in\{1,2\}, i\in\{x,y\}} \left\{ -\frac{\ln(\Lambda_{a,i}(1-P_{a,i}))}{8e} \right\} = \min_{a\in\{1,2\}, i\in\{x,y\}} \left\{ -\frac{\ln(P_{a,i})}{8e} \right\}.
$$

Next, we have that

$$
\mathbb{E}[\exp(\gamma(V^\pi - \mathbb{E}[V^\pi]))] \leq \sum_{a\in\{1,2\}} \mathbb{E}[\exp(\gamma(V^{\pi_a} - \mathbb{E}[V^{\pi_a}])))|type(t) \in \arg\max_{i\in\{1,2\}} P_{a,i}] \cdot \Pr[type(t) \in \arg\max_{i\in\{1,2\}} P_{a,i}] \leq \max_{a\in\{1,2\}} \{ \mathbb{E}[\exp(\gamma(V^{\pi_a} - \mathbb{E}[V^{\pi_a}]))] \}
$$

Where $\hat{V}^{\pi_a}$ is the return for the instance $\langle 1 \rangle, [2], q, \hat{P}, \hat{A} \rangle$ such that for every $a \in \{1,2\}$: $\hat{P}_{a,1} = \max_{i\in\{x,y\}} P_{a,i}$ and $\Lambda_{a,1} = 1$.

Finally, from Lemma 4.3 we get

$$
\max_{a\in\{1,2\}} \{ \mathbb{E}[\exp(\gamma(V^{\pi_a} - \mathbb{E}[V^{\pi_a}]))] \} \leq \max_{a\in\{1,2\}} \exp \left( \frac{8e}{\ln(P_{a,1})} \right)^2 \frac{\gamma^2}{2} = \max_{a\in\{1,2\}, i\in\{x,y\}} \exp \left( \frac{8e}{\ln(P_{a,i})} \right)^2 \frac{\gamma^2}{2}.
$$

Choosing

$$
\tau = b = \max_{a\in\{1,2\}, i\in\{x,y\}} \frac{8e}{\ln(P_{a,i})}
$$

completes the proof as

$$
\max_{a\in\{1,2\}, i\in\{x,y\}} -\frac{8e}{\ln(P_{a,i})} \leq -\frac{8e}{\ln(1-\epsilon)} = \hat{\tau} \quad \text{and} \quad \frac{\tau^2}{b^2} = 1 = \eta.
$$

\textbf{Lemma 5.8.} For every $H \in \mathbb{N}$, it holds that

$$
\mathbb{E} \left[ V^\pi - \max_{\pi \in \Pi_H} V^\pi \right] \leq \frac{1}{2O(H)}.
$$
Proof. Recall that $V^\pi = \sum_{j=1}^{N^\pi} r_j(\pi_j)$, where we drop the dependence on the user index for readability. Formulating differently, for any $H \in \mathbb{N}$ it holds that

$$V^\pi = \sum_{j=1}^{H} \mathbb{I}_{j \leq N^\pi} \cdot r_j(\pi_j) + \sum_{j=H+1}^{\infty} \mathbb{I}_{j \leq N^\pi} \cdot r_j(\pi_j).$$

Using the same representation for $V^{\pi'}$ and taking expectation, we get that

$$\mathbb{E} \left[ V^\pi - V^{\pi'} \right] \leq \mathbb{E} \left[ \sum_{j=1}^{H} \mathbb{I}_{j \leq N^\pi} \cdot r_j(\pi_j) - \sum_{j=1}^{H} \mathbb{I}_{j \leq N^{\pi'}} \cdot r_j(\pi'_j) \right] + \mathbb{E} \left[ \sum_{j=H+1}^{\infty} \mathbb{I}_{j \leq N^\pi} \cdot r_j(\pi_j) \right]$$

$$\leq 0 + \mathbb{E} \left[ \sum_{j=H+1}^{\infty} \mathbb{I}_{j \leq N^\pi} \cdot r_j(\pi_j) \right] = \sum_{j=H+1}^{\infty} \Pr(j \leq N^\pi) r_j(\pi_j)$$

$$\leq \sum_{j=H+1}^{\infty} (1-\epsilon)^j (1-\epsilon) = (1-\epsilon)^{H+2} \sum_{j=0}^{\infty} (1-\epsilon)^j$$

$$\leq (1-\epsilon)^H \frac{1}{\epsilon} \leq \frac{e^{-\epsilon H}}{\epsilon} = \frac{1}{2O(H)}. \quad \square$$