Advanced Quantum Poisson Solver in the NISQ era

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Abstract—The Poisson equation has many applications across the broad areas of science and engineering. Most quantum algorithms for the Poisson solver presented so far, either suffer from lack of accuracy and/or are limited to very small sizes of the problem, and thus have no practical usage. Here we present an advanced quantum algorithm for solving the Poisson equation with high accuracy and dynamically tunable problem size. After converting the Poisson equation to the linear systems through the finite difference method, we adopt the Harrow-Hassidim-Lloyd (HHL) algorithm as the basic framework. Particularly, in this work we present an advanced circuit that ensures the accuracy of the solution by implementing non-truncated eigenvalues through eigenvalue amplification as well as by increasing the accuracy of the controlled rotation angular coefficients, which are the critical factors in the HHL algorithm. We show that our algorithm not only increases the accuracy of the solutions, but also composes more practical and scalable circuits by dynamically controlling problem size in the NISQ devices. We present both simulated and experimental results, and discuss the sources of errors. Finally, we conclude that overall results on the quantum hardware are dominated by the error in the CNOT gates.

Index Terms—Poisson Equation, Quantum Algorithm, Quantum Circuit, HHL Algorithm

I. INTRODUCTION

The Poisson equation is a second order partial differential equations widely used in various fields of science and engineering. In general, in order to solve the Poisson equation numerically, projection methods including such as collocation, spectral, and boundary element methods as well as finite-difference methods [1] are used. The core of these methods is to approximate the solution of the Poisson equation as the solution of linear systems. However, since the dimension of the linear system obtained from the discrete Poisson equation is generally very large, this belongs to the class of problems demanding much computational time.

Therefore, a fascinating technology that would significantly reduce the computational cost of solving the Poisson equation is the application of quantum computing which is faster and more powerful computation than classical computing. Cao et al. [2] firstly used the HHL algorithm [3] to solve the Poisson equation with the quantum circuit model. Later, Wang et al. [4] pointed out a bottleneck of Cao’s algorithm in implementing the controlled rotation of HHL. In their work, they developed new way of implementing the controlled rotation of HHL using named quantum function-value binary expansion (qFBE) [5] and quantum algorithms for solving the reciprocal and square root operations based on the classical non-restoring method. Thus, they not only reduced the algorithm’s complexity, but also made the circuit complete and implementable. However, even though their works apparently improved the quantum algorithm and circuits for the Poisson solver, these still either suffer from lack of accuracy and/or are limited to very small sizes of the problem, and thus have no practical usage.

Here we present an advanced quantum algorithm for solving the Poisson equation with high accuracy and dynamically tunable problem size. Particularly, we develop a circuit based on the Wang’s algorithm that ensures the accuracy of the solution by implementing non-truncated eigenvalues through eigenvalue amplification as well as by increasing the accuracy of the controlled rotation angular coefficients, which are the critical factors in the HHL algorithm.

As a result, we achieve higher success probability which typically decreases with a rate of $1/\kappa^2$, where $\kappa$ is proportional to the number of finite difference discretization, as the problem size is increased. We show that our algorithm not only increases the accuracy of the solutions, but also composes more practical and scalable circuits by dynamically controlling problem size in the NISQ devices. We present both simulated and experimental results, and discuss the sources of errors. Finally, we discuss how the overall results on the currently available quantum hardware are dominated by the error in the CNOT gates.

II. OVERVIEW OF THE PROBLEM

The goal of this work is to implement a quantum algorithm solving the multi-dimensional Poisson equation with boundary conditions. Let us consider the Poisson equation defined in an
open bounded domain $\Omega \subset \mathbb{R}^d$, where $d$ is the number of spatial dimensions.

$$-\nabla^2 v(x) = b(x), \quad x \text{ in } \Omega$$

$$v(x) = 0, \quad x \text{ on } \delta \Omega \quad \Omega = (0,1)^d$$

(1) (2)

where $\delta \Omega$ is the boundary of $\Omega$. One way to solve this problem is to discretize $\Omega$ in $N' = N + 1$ grid points in each dimension, where $N$ is an exponent of base 2 in this work. The solution $v(x)$ is a vector of $(N - 1)^d$ entries.

In this work, we focus on the one-dimensional Poisson equation with Dirichlet boundary conditions. Using the central-difference approximation to discretize the second-order derivative, Eq. (1) could be converted to finite difference form as

$$
A \cdot \begin{pmatrix}
v_1 \\
v_2 \\
\vdots \\
v_{N-1}
\end{pmatrix} = \frac{1}{\pi^2} \begin{pmatrix}
2 & \cdots & -1 & 0 \\
\vdots & \ddots & \vdots & \vdots \\
0 & \cdots & 2 & -1 \\
0 & \cdots & -1 & 2
\end{pmatrix} \begin{pmatrix}
v_1 \\
v_2 \\
\vdots \\
v_{N-1}
\end{pmatrix}
$$

$$A \cdot \begin{pmatrix}
2 & \cdots & -1 & 0 \\
\vdots & \ddots & \vdots & \vdots \\
0 & \cdots & 2 & -1 \\
0 & \cdots & -1 & 2
\end{pmatrix} \begin{pmatrix}
v_1 \\
v_2 \\
\vdots \\
v_{N-1}
\end{pmatrix} = \begin{pmatrix}
b_1 \\
b_2 \\
\vdots \\
b_{N-1}
\end{pmatrix}
$$

(3)

We now have the $N-1$ linear equation system, i.e., $A v = b$ to be solved. Matrix $A$ is a Hermitian matrix and has the dimension of $(N-1) \times (N-1)$ and the mesh size $h$ equals to $1/N$. The best numerical algorithms for solving this problem run polynomially with matrix size [6] in classical computing, so the run-time increases exponentially with the dimension of the problem. In this paper a quantum algorithm is used to produce a quantum state representing the normalized solution of the problem. Since this technique runs in polylog time the curse of dimensionality can be broken. Thus, we will now solve the linear system of equations based on the HHL algorithm [3].

III. QUANTUM ALGORITHM AND CIRCUIT DESIGN

The overall circuit diagram of the present algorithm for solving the one-dimensional Poisson equation is presented in Fig. 1 [4], [7]. As the figure shows, the algorithm consists of the phase estimation, the controlled rotation, and the uncomputation stages. Its circuit diagram has three main registers – reg. B, reg. E, and reg. A. Reg. B is used to encode the coefficients of the right-hand side of Eq. (1) and its number of qubits is $n = \lceil \log(N') \rceil$. Reg. E is used to store the approximated eigenvalues and its number of qubits is $m = 2n + 2 + f$, where the most significant $2n + 2$ bits hold the integer part and the remaining $f$ bits the fractional part. Reg. A is used to store calculated angular coefficients for the controlled rotation operation, and its number of qubits is chosen to be $l \geq m$.

Fig. 1: The overall circuit representation of the algorithm for solving one-dimensional Poisson equation. The number of qubits of registers A, E and B is $l$, $m$ and $n$, respectively. $|\omega_j\rangle$ is the angular coefficient evolved from the approximated eigenvalue $|\lambda_j\rangle$, output of the QPE. The input $|b\rangle_n = \sum_{i=1}^{2^n-1} b_i |i\rangle$ is prepared and stored in register B.

In this work, we assume that the input state $|b\rangle$ of reg. B is prepared as $\sum_i b_i |i\rangle$, where $b_i$ is the value on the right-hand side of Eq.(3), and $|i\rangle$ is the computational basis [8]. That is, the input $|b\rangle$ assigns the pre-requisite state-vector. In other words, reg. B contains the problem that we are trying to solve, which we then entangle with the approximated eigenvalues $\lambda_j$ on the ancillary register. The output of the algorithm thus is a quantum state which encodes the solutions of the Poisson equation as probability amplitudes on reg. B. Thus, this circuit is a process of quantum state preparation, with the output written as: $|v\rangle = A^{-1} |b\rangle = \sum_i \alpha_i |i\rangle$, where $\alpha_i$ is the value of the solutions of the Poisson equation after normalization.

The algorithm used in this work follows several steps:

- Prepare the quantum state $\sum_i b_i |i\rangle$ in reg. B.
- Use Quantum Phase Estimation (QPE) algorithm on regs. B and E. This algorithm applies several Hamiltonian simulations of $U = e^{iAt}$ with $t = 2\pi \frac{1}{\omega_j} 2^k$, $k = 0, ..., n - 1$, to reg. B and entangles the eigenvalues $\lambda_j$ of matrix $A$ in reg. E with the eigenstates $|u_j\rangle$ in reg. B. The system has now the state: $\sum_i \beta_j |\lambda_j\rangle |u_j\rangle$.
- Apply the controlled rotation which consists of two parts: preparing the rotation angular coefficients $|\omega_j\rangle$ in reg. A and performing the controlled $R_\theta$ operation on the ancillary qubit.
- Uncompute QPE and $|\omega_j\rangle$ operations on regs. A, E and B.
- Measure the ancillary qubit. If the measurement of the qubit results in state $|1\rangle$, the algorithm successfully transforms reg. B into the solution $\sum_j \beta_j \frac{1}{\omega_j} |u_j\rangle$. Otherwise the algorithm has to be restarted.

A. Quantum Phase Estimation

Phase estimation is used to approximate the eigenvalues of the discretized matrix $A$ and entangle the states encoding the eigenvalues with the corresponding eigenstates [9]. Hamiltonian simulation of $e^{iAt}$ is the crucial part in phase estimation. Therefore, we first start with exploiting properties of matrix
A to efficiently solve the HHL algorithm by simulating the unitary operator $e^{-iA}$. The eigenvalues of matrix $A$ are $\lambda_j = 4N^2 \sin^2(j \pi / 2N)$. The corresponding eigenvectors are $u_j(k) = \sqrt{2/N} \sin(j \pi k / N)$ [10]. Utilizing the properties of matrix $A$, the unitary operator can be decomposed with a Hermitian matrix $S$ ($S$ being an orthogonal matrix composed of the eigenvectors of $A$), and finally can be diagonalized via the sine transform. After sine transform, phase kickback, and adopt a quantum algorithm for the square root operation based on the classical non-storing method [11]. The detailed circuit composition for phase estimation is discussed in Ref. [4].

B. Controlled Rotation

After phase estimation, we perform the linear map taking the state of $|\lambda_j\rangle$ to $(1/\lambda_j)|\lambda_j\rangle$. This process consists of two parts: calculating the rotation angular coefficients and performing the controlled $R_y$ operation. The probability amplitude of $1/\lambda_j$ would be produced by implementing the controlled $R_y$ rotation, that is, $R_y(2\theta_j)|0\rangle = \cos \theta_j|0\rangle + \sin \theta_j|1\rangle$, where the rotation angle $\theta_j$ can be expressed in terms of $\lambda_j$ as

$$\sin \theta_j = 1/\lambda_j$$

(4)

Which can be rewritten as,

$$\cot \theta_j = \sqrt{\lambda_j^2 - 1}, \quad \theta_j \in (0, \pi/2)$$

(5)

Taking $\theta_j = \omega_j \pi$, then Eq. (5) becomes

$$\omega_j = \frac{1}{\pi} \arccot \left( \sqrt{\lambda_j^2 - 1} \right), \quad \omega_j \in (0, 1/2)$$

(6)

where $\omega_j$ is the rotation angular coefficient. Preparing them through the circuit design has an exponential cost with the problem size [4], so in our implementation we pre-prepare $\omega_j$ and encode them into the circuit.

C. Eigenvalue Amplification

After controlled rotation, the uncomputation is implemented to evolve the state of register B, E, and A back to the initial state. Finally, we perform the measurement operation. If the measurement result of the ancillary qubit is $|1\rangle$, then we get the solution state for the Poisson equation.

However, as the problem size is increased, the success rate of the algorithm is decreased with a rate of $1/\kappa^2$, where $\kappa$ is proportional to the number of finite difference discretization. While expanding the problem size, we find that errors in the eigenvalues start to accumulate in the calculation of the rotation angular coefficient and thus identify a critical source of error in the truncation of the eigenvalues of the matrix $A$.

Therefore, in order to improve the accuracy of the algorithm, we amplify the eigenvalue by a factor of $2^i$ ($i$ is an integer) in QPE and making sure that it is being compensated by a normalization factor in the later section of the circuit [7]. This amplification is critical to getting accurate results from the larger matrices, as the success rate was so low that the number of times required to run the circuit quickly becomes a limiting factor. However, the eigenvalue amplification significantly improves the accuracy of the of the computation resulting in a higher success probability.

IV. RESULTS AND CONCLUSIONS

We first demonstrate our algorithm and circuit with IBM’s QASM and MPS simulators and produce the result for relatively a smaller problem defined by the size of the matrix $A$ in Eq. (3). This is shown in Fig. 2, that compares the results of a $4 \times 4$ problem with $|b\rangle = 0.0|00\rangle + \frac{1}{\sqrt{2}}|01\rangle + \frac{1}{2}(|10\rangle + |11\rangle)$ being the right-hand side of the Poisson equation. Though there is some discrepancies of QASM-based result, the MPS-based solution shows reasonable agreement with both analytical and existing QRUNES [4] results. The accuracy of the MPS-based result is depicted through the relative error in MPS solution with respect to the analytical result, shown in the right side of Fig. 2.

The circuit of $4 \times 4$ size of the problem uses 32 qubits which is exactly the limit of QASM simulator. In contrast, MPS simulator allows up to 100 qubits. Therefore, as a viable option for presenting the circuit of larger problem sizes, we decided to proceed with the MPS simulator only. The result for a $8 \times 8$ problem with $|b\rangle = 0.0|00\rangle + \frac{1}{\sqrt{2}}|01\rangle + \frac{1}{2}(|10\rangle
Fig. 4: For a $4 \times 4$ problem, comparison of the experimental results, pursued on the *ibm_brooklyn* quantum device, with MPS-based solution. Experiment was done on optimized circuits compiled with different levels of transpiler optimization parameters and sampling. The cross indicates an experimental artifact due to the hardware error involving CNOT gates.

\[
|011\rangle + |100\rangle + \frac{1}{2}(|101\rangle + |110\rangle + |111\rangle)
\]

is shown in Fig. 3. With the implementation of the eigenvalue amplification by a factor of $2^4$ has reduced the relative error by about half, and we are confident that a higher amplification factor (for example, $2^8$) would further reduce this error. The right part of this figure explicitly shows the relative error at the level of the quantum states.

Finally, for a $4 \times 4$ problem, we experiment our circuit on the IBM Brooklyn 65 qubit system (Hummingbird r2) and present the results along with MPS-based solution in Fig. 4. The circuit compiled with different level of transpiler optimization parameter and sampling [12] are used, yet the experimental results come out as erroneous with an artifact of nonzero contribution for $|00\rangle$ state. The average CNOT error on Brooklyn system is $8.094e-2$, in other words they have an accuracy of about 0.99. Thus, we can estimate the total accuracy of the experiment based on the final number of CNOT gates transpiled from the more abstract circuit, approximated by 0.99$^c$ where $c$ is the final number of CNOT gates after transpilation. Our optimized circuit has about 5.5k CNOT gate and their accumulated error results in washing out the accuracy of the overall experimental solution.

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