ABSTRACT
This paper introduces a new metaobject, the generalizer, which complements the existing specializer metaobject. With the help of examples, we show that this metaobject allows for the efficient implementation of complex non-class-based dispatch within the framework of existing metaobject protocols. We present our modifications to the generic function invocation protocol from the *Art of the Metaobject Protocol*; in combination with previous work, this produces a fully-functional extension of the existing mechanism for method selection and combination, including support for method combination completely independent from method selection. We discuss our implementation, within the SBCL implementation of Common Lisp, and in that context compare the performance of the new protocol with the standard one, demonstrating that the new protocol can be tolerably efficient.

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1. INTRODUCTION
The revisions to the original Common Lisp language [16] included the detailed specification of an object system, known as the Common Lisp Object System (CLOS), which was eventually standardized as part of the ANSI Common Lisp standard [12]. The object system as presented to the standardization committee was formed of three chapters. The first two chapters covered programmer interface concepts and the functions in the programmer interface [17, Chapter 28] and were largely incorporated into the final standard; the third chapter, covering a Metaobject Protocol (MOP), was not.

Nevertheless, the CLOS MOP continued to be developed, and the version documented in [8] has proven to be a reasonably robust design. While many implementations have derived their implementations of CLOS from either the Cloette illustrative implementation in [8], or the Portable Common Loops implementation of CLOS from Xerox Parc, there have been largely from-scratch reimplementations of CLOS (in CLISP† and CCL‡, at least) incorporating substantial fractions of the Metaobject Protocol as described.

Although it has stood the test of time, the CLOS MOP is neither without issues (e.g., semantic problems with make-method-lambda [1]; useful functions such as compute-effective-slot-definition-initargs being missing from the standard) nor is it a complete framework for the metaprogrammer to implement all conceivable variations of object-oriented behaviour. While metaprogramming offers

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1 GNU CLISP, at [http://www.clisp.org/](http://www.clisp.org/)
2 Clozure Common Lisp, at [http://ccl.clozure.com/](http://ccl.clozure.com/)
some possibilities for customization of the object system behaviour, those possibilities cannot extend arbitrarily in all directions (conceptually, if a given object system is a point in design space, then a MOP for that object system allows exploration of a region of design space around that point; see figure 1). In the case of the CLOS MOP, there is still an expectation that functionality is implemented with methods on generic functions, acting on objects with slots; it is not possible, for example, to transparently implement support for “message not understood” as in the message-passing paradigm, because the analogue of messages (generic functions) need to be defined before they are used.

Nevertheless, the MOP is flexible, and is used for a number of things, including: documentation generation (where introspection in the MOP is used to extract information from a running system[3]; object-relational mapping[1] and other approaches to object persistence[11]; alternative backing stores for slots (hash-tables[7] or symbols[3]); and programmatic construction of metaobjects, for example for interoperability with other language runtimes' object systems.

One area of functionality where there is scope for customization by the metaprogrammer is in the mechanics and semantics of method applicability and dispatch. While in principle AMOP allows customization of dispatch in various different ways (the metaprocessor can define methods on protocol functions such as compute-applicable-methods, compute-applicable-methods-using-classes), for example, in practice implementation support for this was weak until relatively recently[6].

Another potential mechanism for customizing dispatch is implicit in the class structure defined by AMOP: standard specializer objects (instances of class and eql-specializer) are generalized instances of the specializer protocol class, and in principle there are no restrictions on the metaprogrammer constructing additional subclasses. Previous work[10] has explored the potential for customizing generic function dispatch using extended specializers, but the metaprogrammer must override the entirety of the generic function invocation protocol (from compute-discriminating-function on down), leading to toy implementations and duplicated effort.

This paper introduces a protocol for efficient and controlled handling of new subclasses of specializer. In particular, it introduces the generalizer protocol class, which generalizes the return value of class-of in method applicability computation, and allows the metaprogrammer to hook into caching schemes to avoid needless recomputation of effective methods for sufficiently similar generic function arguments (See Figure 2).

The remaining sections in this paper can be read in any order. We give some motivating examples in section 2 including reimplementations of examples from previous work, as well as examples which are poorly supported by previous protocols. We describe the protocol itself in section 3.

As in many of the systems surveyed at https://sites.google.com/site/sabraonthehill/lisp-document-generation-apps[4], the Closer to MOP project at http://common-lisp.net/project/closer/[5] attempts to harmonize the different implementations of the metaobject protocol in Common Lisp.

Figure 2: Dispatch Comparison

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2. EXAMPLES

In this section, we present a number of examples of dispatch implemented using our protocol, which we describe in section 3. For reasons of space, the metaprocessor code examples in this section do not include some of the necessary support code to run; complete implementations of each of these cases, along with the integration of this protocol into the SBCL implementation[13] of Common Lisp, are included in the authors' repository[6].

A note on terminology: we will attempt to distinguish between the user of an individual case of generalized dispatch (the "programmer"), the implementor of a particular case of generalized dispatch (the "metaprogrammer"), and the authors as the designers and implementors of our generalized dispatch protocol (the "metametaprogrammer", or more likely “we”).

2.1 CONS specializers

One motivation for the use of generalized dispatch is in an extensible code walker: a new special form can be handled simply by writing an additional method on the walking generic function, seamlessly interoperating with all existing methods. In this use-case, dispatch is performed on the first element of lists. Semantically, we allow the programmer to specialize any argument of methods with a new kind of specializer, cons-specializer, which is applicable if and only if the corresponding object is a cons whose car is eql to the symbol associated with the cons-specializer; these specializers are more specific than the cons class, but less specific than an eql-specializer on any given cons.

The programmer code using these specializers is unchanged from[10]; the benefits of the protocol described here are: that the separation of concerns is complete – method selection is independent of method combination – and that...
the protocol allows for efficient implementation where possible, even when method selection is customized. In an application such as walking source code, we would expect to encounter special forms (distinguished by particular atoms in the car position) multiple times, and hence to dispatch to the same effective method repeatedly. We discuss the efficiency aspects of the protocol in more detail in section 3.1.2; we present the metaprogrammer code to implement the cons-specializer below.

(defclass cons-specializer (specializer)
  ((%car :reader %car :initarg :car))
(defclass cons-generalizer (generalizer)
  ((%car :reader %car :initarg :car)))
(defmethod generalizer-of-using-class
  ((gf cons-generic-function) arg)
  (typecase arg
    ((cons symbol)
      (make-instance 'cons-generalizer :car (car arg))))
    (t (call-next-method))))
(defmethod generalizer-equal-hash-key
  ((gf cons-generic-function) (arg real))
  (%car g))
(defmethod specializer-accepts-generalizer-p
  ((gf cons-generic-function) (s cons-specializer)
   (g cons-generalizer))
  (if (eql (%car s) (%car g))
    (values t t)
    (values nil t)))
(defmethod specializer-accepts-p
  ((s cons-specializer) o)
  (and (consp o) (eql (car o) (%car s))))

The code above shows a minimal use of our protocol. We have elided some support code for parsing and unparsing specializers, and for handling introspective functions such as finding generic functions for a given specializer. We have also elided methods on the protocol functions specializer< and same-specializer-p: for cons-specializer objects, specializer ordering is trivial, as only one cons-specializer (up to equality) can ever be applicable to any given argument. See section 2.3 for a case where specializer ordering is non-trivial.

As in [10], the programmer can use these specializers to implement a modular code walker, where they define one method per special operator. We show two of those methods below, in the context of a walker which checks for unused bindings and uses of unbound variables.

(defgeneric walk (form env stack)
  (:generic-function-class cons-generic-function))
(defmethod walk
  ((expr (cons lambda)) env call-stack)
  (let (((lambda-list (cadr expr))
           (body (cddr expr)))
         (with-checked-bindings
           ((bindings-from-l1 lambda-list)
            env call-stack))
         (dolist form body)
        form)
  (walk form env (cons form call-stack))))

Note that in this example there is no strict need for cons-specializer and cons-generalizer to be distinct classes. In standard generic function dispatch, the class functions both as the specializer for methods and as the generalizer for generic function arguments; we can think of the dispatch implemented by cons-specializer objects as providing for subclasses of the cons class distinguished by the car of the cons. This analogy also characterizes those use cases where the metaprogrammer could straightforwardly use filtered dispatch [2] to implement their dispatch semantics. We will see in section 2.3 an example of a case where filtered dispatch is incapable of straightforwardly expressing the dispatch, but first we present our implementation of the motivating case from [2].

2.2 SIGNUM specializers

Our second example of the implementation and use of generalized specializers is a reimplementation of one of the examples in [2]: specifically, the factorial function. Here, dispatch will be performed based on the signum of the argument, and again, at most one method with a signum specializer will be applicable to any given argument, which makes the structure of the specializer implementation very similar to the cons specializers in the previous section.

The metaprogrammer has chosen in the example below to compare signum values using =, which means that a method with specializer (signum 1) will be applicable to positive floating-point arguments (see the first method on specializer-accepts-generalizer-p and the method on specializer-accepts-p below). This leads to one subtle difference in behaviour compared to that of the cons specializers: in the case of signum specializers, the next method after any signum specializer can be different, depending on the class of the argument. This aspect of the dispatch is handled by the second method on specializer-accepts-generalizer-p below.

(defclass signum-specializer (specializer)
  (%signum :reader %signum :initarg :signum))
(defclass signum-generalizer (generalizer)
  (%signum :reader %signum :initarg :signum))
(defmethod generalizer-of-using-class
  ((gf signum-generic-function) arg real)
  (make-instance 'signum-generalizer :signum (signum arg)))
(defmethod generalizer-equal-hash-key
  ((gf signum-generic-function) (arg real))
  (make-instance 'signum-generalizer :signum (signum arg)))

Given these definitions, and once again some more straightforward ones elided for reasons of space, the programmer can implement the factorial function as follows:

```
(defun fact (n)
  (if (<= n 1) 1
      (* n (fact (1- n))))
```

The programmer does not need to include a method on \(\text{signum} -1\), as the standard `no-applicable-method` protocol will automatically apply to negative real or non-real arguments.

### 2.3 Accept HTTP header specializers

In this section, we implement a non-trivial form of dispatch. The application in question is a web server, and specifically to allow the programmer to support RFC 2616 \[5\] content negotiation, of particular interest to publishers and consumers of REST-style Web APIs.

The basic mechanism in content negotiation is as follows: the web client sends an HTTP request with an `Accept` header, which is a string describing the media types it is willing to receive as a response to the request, along with numerical preferences. The web server compares these stated client preferences with the resources it has available to satisfy this request, and sends the best matching resource in its response.

For example, a graphical web browser might send an `Accept` header of\[3\]

```
Accept: text/html, application/xml;q=0.9, */*;q=0.8
```

for a request of a resource typed in to the URL bar. This should be interpreted as meaning that: if the server can provide content of type `text/html` (i.e. HTML) for that resource, then it should do so. Otherwise, if it can provide `application/xml` content (i.e. XML of any schema), then that should be provided; failing that, any other content type is acceptable.

In the case where there are static files on the filesystem, and the web server must merely select between them, there is not much more to say. However, it is not unusual for a web service to be backed by some other form of data, and responses computed and sent on the fly, and in these circumstances the web server must compute which of its known output formats it can use to satisfy the request before actually generating the best matching response. This can be modelled as one generic function responsible for generating the response, with methods corresponding to content-types – and the generic function must then perform method selection against the request’s `Accept` header to compute the appropriate response.

The `accept-specializer` below implements this dispatch. It depends on a lazily-computed `tree` slot to represent the information in the accept header (generated by `parse-accept-string`), and a function \(q\) to compute the (defaulted) preference level for a given content-type and `tree`; then, method selection and ordering involves finding the \(q\) for each `accept-specializer`’s content type given the `tree`, and sorting them according to the preference level.

```
(defclass accept-specializer (specializer)
  ((media-type :initarg :media-type :reader media-type)))

(defclass accept-generalizer (generalizer)
  ((header :initarg :header :reader header)
   (tree)
   (next :initarg :next :reader next)))

(defmethod generalizer-equal-hash-key
  ((gf accept-generic-function) (g accept-generalizer))
  '(:accept-generalizer ,((header g))))

(defmethod specializer-accepts-generalizer-p
  ((gf accept-generic-function) (s accept-specializer) (g accept-generalizer))
  (values (q (media-type s) (tree g)) t))

(defmethod specializer-accepts-generalizer-p
  ((gf accept-generic-function) (s specializer) (g accept-generalizer))
  (specializer-accepts-generalizer-p gf s (next g)))

(defmethod specializer-<
  ((gf accept-generic-function) (s1 accept-specializer) (s2 accept-specializer) (g accept-generalizer))
  (let ((m1 (media-type s1)) (m2 (media-type s2)))
    (tree (tree g)))
  (cond
    ((string= m1 m2) '=)
    (t (let ((q1 (q m1 tree)) (q2 (q m2 tree)))
        (cond
          ((= q1 q2) '=)
          (((< q1 q2) '>)
            (t '<)))))))
```

The metaprogrammer can then add support for objects representing client requests, such as instances of the `request` class in the Hunchentoot \[7\] web server, by translating these

\[7\] Hunchentoot is a web server written in Common Lisp, al-
into accept-generalizer instances. The code below implements this, by defining the computation of a generalizer object for a given request, and specifying how to compute whether the specializer accepts the given request object (q returns a number between 0 and 1 if any pattern in the tree matches the media type, and nil if the media type cannot be matched at all).

```
(let* ((accept (tbnl:header-in :accept o))
   (arg tbnl:request))
  (and q (> q 0)))
```

This dispatch cannot be implemented using filtered dispatch, except by generating anonymous classes with all the right mime-types as direct superclasses in dispatch order; the filter would generate

```
(ensure-class nil :direct-superclasses
  `(text/html image/webp ...))
```

and dispatch would operate using those anonymous classes. While this is possible to do, it is awkward to express content-type negotiation in this way, as it means that the dispatcher must know about the universe of mime-types that clients might declare that they accept, rather than merely the set of mime-types that a particular generic function is capable of serving; handling wildcards in accept strings is particularly awkward in the filtering paradigm.

Note that in this example, the method on specializer< involves a non-trivial ordering of methods based on the q values specified in the accept header (whereas in sections 2.1 and 2.2 only a single extended specializer could be applicable to any given argument).

Also note that the accept specializer protocol is straightforwardly extensible to other suitable objects; for example, one simple debugging aid is to define that an accept-specializer should be applicable to string objects. This can be done in a modular fashion (see the code below, which can be completely disconnected from the code for Hunchentoot request objects), and generalizes to dealing with multiple web server libraries, so that content-negotiation methods are applicable to each web server’s request objects.

```
(defmethod specializer-accepts-p
  ((s accept-specializer) (o string))
  (and q (> q 0)))
```

The next slot in the accept-generalizer is used to deal with the case of methods specialized on the classes of objects as well as on the acceptable media types; there is a method on specializer-accepts-generalizer< for specializers that are not of type accept-specializer which calls the generic function again with the next generalizer, so that methods specialized on the classes of objects and string are treated as applicable to corresponding objects, though less specific than methods with accept-specializer specializations.

3. PROTOCOL
In section 2 we have seen a number of code fragments as partial implementations of particular non-standard method dispatch strategies, using generalizer metaobjects to mediate between the methods of the generic function and the actual arguments passed to it. In section 3.1 we go into more detail regarding these generalizer metaobjects, describing the generic function invocation protocol in full, and showing how this protocol allows a similar form of effective method caching as the standard one does. In section 3.2 we show the results of some simple performance measurements on our implementation of this protocol in the SBCL implementation [13] of Common Lisp to highlight the improvement that this protocol can bring over a naive implementation of generalized dispatch, as well as to make the potential for further improvement clear.

3.1 Generalizer metaobjects
3.1.1 Generic function invocation
As in the standard generic function invocation protocol, the generic function’s actual functionality is provided by a discriminating function. The functionality described in this protocol is implemented by having a distinct subclass of standard-generic-function, and a method on compute-discriminating-function which produces a custom discriminating function. The basic outline of the discriminating function is the same as the standard one: it must first compute the set of applicable methods given particular arguments; from that, it must compute the effective method by combining the methods appropriately according to the generic function’s method combination; finally, it must call the effective method with the arguments.

Computing the set of applicable methods is done using a pair of functions: compute-applicable-methods, the standard metaobject function, and a new function compute-applicable-methods-using-generalizers. We define a custom method on compute-applicable-methods which tests the applicability of a particular specializer against a given argument using specializer-accepts-p, a new protocol function with default implementations on class and eql-specializer to implement the expected behaviour. To order the methods, as required by the protocol, we define a pairwise comparison operator specializer< which defines an ordering between specializers for a given generalizer argument (remembering that even in standard CLOS the or-
dering between class specializers can change depending on
the actual class of the argument).

The new compute-applicable-methods-using-generalizers
is the analogue of the MOP’s compute-applicable-methods-
using-classes. Instead of calling it with the class-of each
argument, we compute the generalizers of each argument us-
ing the new function generalizer-of-using-class (where the
method rather than the class of the object), and call compute-
applicable-methods-using-generalizers with the generic
function and list of generalizers. As with compute-applicable-
methods-using-classes, a secondary return value indicates
whether the result of the function is definitive for that list of
generalizers.

Thus, in generic function invocation, we first compute the
generalizers of the arguments; we compute the ordered set of
applicable methods, either from the generalizers or (if that
is not definitive) from the arguments themselves; then the
normal effective method computation and call can occur.
Unfortunately, the nature of an effective method function is
not specified, so we have to reach into implementation
internals a little in order to call it, but otherwise the remain-
der of the generic function invocation protocol is unchanged
from the standard one. In particular, method combination
is completely unchanged; programmers can choose arbitrary
method combinations, including user-defined long form com-
binations, for their generic functions involving generalized
dispatch.

3.1.2 Effective method memoization
The potential efficiency benefit to having generalizer
metaobjects lies in the use of compute-applicable-methods-
using-generalizers. If a particular generalized special-
izer accepts a variety of objects (such as the signum special-
izer accepting all reals with a given sign, or the ac-
cept specializer accepting all HTTP requests with a par-
ticular Accept header), then there is the possibility of
caching and reusing the results of the applicable and effec-
tive method computation. If the computation of the applic-
able method from compute-applicable-methods-using-
generalizers is definitive, then the ordered set of applicable
methods and the effective method can be cached.

One issue is what to use as the key for that cache. We cannot
use the generalizers themselves, as two generalizers that
should be considered equal for cache lookup will not com-
pare as equal – and indeed even the standard generalizer,
the class, cannot easily be used as we must be able to in-
validade cache entries upon class redefinition. The issue of
class generalizers we can solve as in [9] by using the wrapper
of a class, which is distinct for each distinct (re)definition of
a class; for arbitrary generalizers, however, there is a priori
no good way of computing a suitable hash key automatically,
as we allow the metaprogrammer to specify one by defining
a method on generalizer-equal-hash-key, and combining
the hash keys for all required arguments in a list to use as a
key in an equal hash-table.

3.2 Performance
We have argued that the protocol presented here allows for
expressive control of method dispatch while preserving the
possibility of efficiency. In this section, we quantify the ef-
ficiency that the memoization protocol described in section
The relative picture is similar to the `cons-specializer` case; including a cache saves a factor of 10 in this case, and another factor of 3 for the one-argument cache special case. The cost of the genericity of the protocol here is starker; even the one-argument cache is a factor of 6 slower than the standard generic-function implementation, and a further factor of 2 away from the implementation of factorial as a function. We discuss ways in which we expect to be able to improve performance in section 5.1.

We could allow the metaprogrammer to improve on the one-argument performance by constructing a specialized cache: for `signum` arguments of rational arguments, the logical cache structure is to index a three-element vector with (1+ `signum`). The current protocol does not provide a way of eliding the two generic function calls for the generic cache; we discuss possible approaches in section 6.

3.3 Full protocol

The protocol described in this paper is only part of a complete protocol for specializer and generalizer metaobjects. Our development of this protocol is as yet incomplete; the work described here augments that in [10], but is yet relatively untested – and additionally our recent experience of working with that earlier protocol suggests that there might be useful additions to the handling of specializer metaobjects, independent of the generalizer idea presented here.

4. RELATED WORK

The work presented here builds on specializer-oriented programming described in [10]. Approximately contemporaneously, filtered dispatch [2] was introduced to address some of the same use cases; filtered dispatch works by having a custom discriminating function which wraps the usual one, where the wrapping function augments the set of applicable methods with applicable methods from other (hidden) generic functions, one per filter group; this step is not memoized, and using eql methods to capture behaviours of equivalence classes means that it is hard to see how it could be. The methods are then combined using a custom method combination to mimic the standard one; in principle implementors of other method combinations could cater for filtered dispatch, but they would have to explicitly modify their method combinations. The Clojure programming language supports multimethods [8] with a variant of filtered dispatch as well as hierarchical and identity-based method selectors.

In context-oriented programming [6, 13], context dispatch occurs by maintaining the context state as an anonymous class with the superclasses representing all the currently active layers; this is then passed as a hidden argument to context-aware functions. The set of layers is known and under programmer control, as layers must be defined beforehand.

In some sense, all dispatch schemes are specializations of predicate dispatch [4]. The main problem with predicate dispatch is its expressiveness: with arbitrary predicates able to control dispatch, it is essentially impossible to perform any substantial precomputation, or even to automatically determine an ordering of methods given a set of arguments.

Even Clojure’s restricted dispatch scheme provides an explicit operator for stating a preference order among methods, where here we provide an operator to order specializers; in filtered dispatch the programmer implicitly gives the system an order of precedence, through the lexical ordering of filter specification in a filtered function definition.

The Slate programming environment combines prototype-oriented programming with multiple dispatch [15]; in that context, the analogue of an argument’s class (in Common Lisp) as a representation of the equivalence class of objects with the same behaviour is the tuple of roles and delegations: objects with the same roles and delegations tuple behave the same, much as objects with the same generalizer have the same behaviour in the protocol described in this paper.

The idea of generalization is of course not new, and arises in other contexts. Perhaps of particular interest is generalization in the context of partial evaluation; for example, [14] considers generalization in online partial evaluation, where sets of possible values are represented by a type system construct representing an upper bound. Exploring the relationship between generalizer metaobjects and approximation in type systems might yield strategies for automatically computing suitable generalizers and cache functions for a variety of forms of generalized dispatch.

5. CONCLUSIONS

In this paper, we have presented a new generalizer metaobject protocol allowing the metaprogrammer to implement in a straightforward manner metaobjects to implement custom method selection, rather than the standard method selection as standardized in Common Lisp. This protocol seamlessly interoperates with the rest of CLOS and Common Lisp in general; the programmer (the user of the custom specializer metaobjects) may without constraints use arbitrary method combination, intercede in effective method combination, or write custom method function implementations. The protocol is expressive, in that it handles forms of dispatch not possible in more restricted dispatch systems, while not suffering from the indeterminism present in predicate dispatch through the use of explicit ordering predicates.

The protocol is also reasonably efficient; the metaprogrammer can indicate that a particular effective method computation can be memoized, and under those circumstances much of the overhead is amortized (though there remains a substantial overhead compared with standard generic-function or regular function calls). We discuss how the efficiency could be improved below.

5.1 Future work

Although the protocol described in this paper allows for a more efficient implementation, as described in section 5.1.2 than computing the applicable and effective methods at each generic function call, the efficiency is still some way away from a baseline of the standard generic-function, let alone a standard function. Most of the invocation protocol is memoized, but there are still two full standard generic-function calls – `generalizer-of-using-class` and `generalizer-equal-hash-key` – per argument per call to a generic function with extended specializers, not to mention a hash table lookup.

For many applications, the additional flexibility afforded by
generalized specializers might be worth the cost in efficiency, but it would still be worth investigating how much the overhead from generalized specializers can be reduced; one possible avenue for investigation is giving greater control over the caching strategy to the metaprogrammer.

As an example, consider the \texttt{signum-specializer}. The natural cache structure for a single argument generic function specializing on \texttt{signum} is probably a four-element vector, where the first three elements hold the effective methods for \texttt{signum} values of -1, 0, and 1, and the fourth holds the cached effective methods for everything else. This would make the invocation of such functions very fast for the (presumed) common case where the argument is in fact a real number. We hope to develop and show the effectiveness of an appropriate protocol to allow the metaprogrammer to construct and exploit such caching strategies, and (more speculatively) to implement the lookup of an effective method function in other ways.

We also aim to demonstrate support within this protocol for some particular cases of generalized specializers which seem to have widespread demand (in as much as any language extension can be said to be in “demand”). In particular, we have preliminary work towards supporting efficient dispatch over pattern specializers such as implemented in the \texttt{Optima} library\cite{10} and over a prototype object system similar to that in \texttt{Slate}\cite{14}. Our current source code for the work described in this paper can be seen in the git source code repository at \url{http://christophe.rhodes.io/git/specializable.git} which will be updated with future developments.

Finally, after further experimentation (and, ideally, non-trivial use in production) if this protocol stands up to use as we hope, we aim to produce a standards-quality document so that other implementors of Common Lisp can, if they choose, independently reimplement the protocol, and so that users can use the protocol with confidence that the semantics will not change in a backwards-incompatible fashion.

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