SLIME MOULD COMPUTES PLANAR SHAPES

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ABSTRACT. Computing a polygon defining a set of planar points is a classical problem of modern computational geometry. In laboratory experiments we demonstrate that a concave hull, a connected $\alpha$-shape without holes, of a finite planar set is approximated by slime mould Physarum polycephalum. We represent planar points with sources of long-distance attractants and short-distance repellents and inoculate a piece of plasmodium outside the data set. The plasmodium moves towards the data and envelops it by pronounced protoplasmatic tubes.

Keywords: Unconventional computing, Physarum polycephalum, Concave hull, Concave hull, Alpha shape

1. INTRODUCTION

A slime mould Physarum polycephalum has a rich life cycle [30], which includes fruit bodies, spores, single-cell amoebas, and syncytium. Plasmodium is a vegetative stage of Physarum polycephalum, syncytium, a single cell with many nuclei. The plasmodium consumes microscopic particles. During its foraging behaviour the plasmodium spans scattered sources of nutrients with a network of protoplasmic tubes. The protoplasmic network is optimised to cover all sources of food and to deliver robust and speedy transportation of nutrients and metabolites in the plasmodium body. Plasmodium’s foraging behaviour can be interpreted as computation, when data are represented by spatial configurations of attractants and repellents, and results by structures of protoplasmic network [3, 2]. Plasmodium satisfactory solves many computational problems with natural parallelism, including shortest path [23, 24], implementation of storage modification machines [4], Voronoi diagram [29], logical computing [33, 5], process algebra [28]; see overview in [2].

A convex hull of a finite planar set is a smallest region of the plane containing the set [27]. Classical algorithms of convex hull construction include Jarvis’ gift wrapping [17], Preparata-Hong’s divide and conquer [26], Graham’s scan of pre-sorted set of points [19], and Akl-Toussaint’s quick algorithm based on removing points lying inside a convex quadrilateral of extreme points [7]: see reviews and discussion of the algorithms in [27, 11] and codes of implementations in [25]. Multi-processor and systolic-processor solutions of the convex hull problem are mostly based on parallelisation of the classical serial algorithms [10] [15] or exploring extremal properties of the given set [18, 21, 12], see also review in [8]. Cellular-automata algorithm are based on propagating patterns halting their growth in the hull’s boundaries [1] or encapsulating the data set and developing synchronisation patterns [32, 9], are theoretical precursors for bio-inspired approach to computation of convex hulls.

A convex hull is economical but rough reconstruction of an object shape from the object’s sample points. It often lacks details necessary for realistic visualisation of objects. In 1983 Edelsbrunner, Kirkpatrick and Seidel [13] introduced $\alpha$-shapes, generalisation of convex hulls, where control parameter $\alpha$ allows for a smooth transition between crude...
and fine approximations of point sets. There are hundreds of algorithms for computing $\alpha$-shapes yet no experimental laboratory prototypes of chemical, physical or biological computers for approximation of $\alpha$-shapes. Present paper partly fills the gap and shows how to approximate $\alpha$-shapes without holes using foraging behaviour of $P.\ polycephalum$.

2. Experimental

A convex hull of a finite set $P$ of planar points is the smallest convex polygon that contains all points of $P$ (Fig. 1a). $\alpha$-hull of $P$ is an intersection of the complement of all closed discs of radius $1/\alpha$ that includes no points of $P$ [13, 14]. $\alpha$-shape is a convex hull when $\alpha \rightarrow \infty$. With decrease of $\alpha$ the shapes may shrink, develop holes and become disconnected, the shapes collapse to $P$ when $\alpha \rightarrow 0$. A concave hull is non-convex polygon representing area occupied by $P$. A concave hull is a connected $\alpha$-shape without holes (Fig. 1b).

**Problem 1.** Given planar set $P$ represented by physical objects plasmodium of $P.\ polycephalum$ must represent concave hull of $P$ by its largest protoplasmic tube.

We cultivate $P.\ polycephalum$ in plastic containers, on paper kitchen towels sprinkled with still drinking water and fed with oat flakes. For experiments we use polystyrene Petri dishes (round, diameter 120 mm, and rectangular 120 $\times$ 120 mm) and 2% agar gel (Select agar, Sigma Aldrich) as a non-nutrient substrate. Images of plasmodium are recorded by scanning Petri dishes in Epson Perfection 4490. Photos are taken using FujiPix 6000 camera.

3. Results

The fist algorithm of convex hull construction [17] was based on cognitive tactic techniques we use in our everyday’s life. We select a starting point which is extremal point of $P$. We pull a rope (anti-)clockwise to other extremal point. We continue until the set $P$ is wrapped completely. The computation stops when we reach the starting point. Let we represent data points $P$ by sources of attractants only, e.g. by oat flakes (Fig. 1cde). We place a piece of plasmodium at some distance away from the set of points, see scheme in Fig. 1e. The plasmodium propagates towards set $P$, colonises oat flakes (Fig. 1cd) and spans them with a network of protoplasmic tubes (Fig. 1e). No hull is constructed. When all data points $P$ are colonised and spanned by protoplasmic network the plasmodium ventures to explore the space around $P$.

**Proposition 1.** $P.\ polycephalum$ does not compute concave or convex hull of a set represented by attracting sources.

Using repellents only, e.g. as a set of obstacles between attractants and inoculation site, will do no good because plasmodium will just pass around the repellents and leave (see details in [2], controlling Physarum with salt). The only solution would be to employ attractants to ‘pull’ plasmodium towards planar set $P$ and to use repellents to prevent plasmodium from spanning the points of $P$ (Fig. 2a). Strength of repellents should be proportional to $\alpha$ and thus will determine exact shape of the constructed hull. This corresponds to original definition [13] that $\alpha$-hull of $P$ is the intersection of all closed discs with radius $1/\alpha$ that contain all the points of $P$.

**Proposition 2.** Plasmodium of $P.\ polycephalum$ approximates connected $\alpha$-hull (without holes) of a finite planar set, which points are represented by sources of long-distance attractants and short-distance repellents.
Figure 1. Basics. (a) Example of convex hull. (b) Example of concave hull. Points of set $P$ are empty discs. (c–d) Spanning oat flakes, representing $P$, by a network of protoplastic tubes: scanned image (c) and its binarisation (d). (e) Scheme of the plasmodium propagation, circles are points, or oat flakes, of $P$ and star marks site of inoculation, arrows are protoplastic tubes.
**Figure 2.** Experimental. (a) Proposed distribution of attracting and repelling gradients which may force plasmodium to approximate a concave hull (arrows aiming towards set of discs are attractive forces, arrows originating in data points, discs, are repelling forces). (bc) Plasmodium’s interaction with the single half-pill of Kalms Tablets: scanned (b) and binarised (c) images; at the beginning of experiment plasmodium is inoculated in southmost part of Petri dish (pill of Kalms Tablets is at the centre of the Petri dish). (d) Typical experimental setup, pills representing \( P \) and the initial site of plasmodium’s inoculation are shown by arrows.
To prove the proposition constructively we must find an appropriate representation of \( P \) and demonstrate in experiments viability of the approach. While experimenting with different potential candidates for combined representation of attractants and repellents we found that the plasmodium’s reaction to pills of Kalms Tablets and Kalms Sleep\(^1\) is somewhat unusual. When presented with a half-pill of the Kalms Tablets/Sleep the plasmodium propagates towards the pill and forms, with its protoplasmic tubes, a circular enclosure around the pill (Fig. 2bc). Such a unique behaviour of plasmodium in presence of Kalms Tablets/Sleep indicates that a plasmodium could implement Jarvis’s Gift Wrapping algorithm \([17]\), adapted to concave hulls without holes, if points of \( P \) are represented by the pills.

We tested feasibility of the idea in 25 experiments. All experiments were successful. In each experiment we arranged 4-8 half-pills (they represented given planar set \( P \)) in a random fashion near centre of a Petri dish and inoculated an oat flake colonised by plasmodium 2-4 cm away from the set \( P \). A typical experiment is illustrated in Fig. 3a–d. In 12 h after inoculation plasmodium propagates towards set \( P \) and starts enveloping the set with its body and network of protoplasmic tubes (Fig. 3b). The plasmodium completes approximation of a shape by entirely enveloping \( P \) in next 12 h (Fig. 3d). The plasmodium does not propagate inside configuration of pills (Fig. 3e).

Configuration of \( P \) in Fig. 3a–d favours approximation of a convex hull. If spatial configuration of points curves inwards then concave hull is approximated (Fig. 4). Typical hulls approximated by plasmodium of \( P. polycephalum \) in our experiments are shown in Fig. 5. Hulls constructed by plasmodium match their counterparts calculated by classical algorithms.

Let us verify experimental laboratory results with computer simulation. A profile of plasmodium’s propagating zone is isomorphic to shapes of wave-fragments in sub-excitable media \([2]\). When active zone of \( P. polycephalum \) propagates two processes occur simultaneously — propagation of the wave-shaped tip of the pseudopodium and formation of the trail of protoplasmic tubes. We simulate the (chemo-)tactic traveling of plasmodium using two-variable Oregonator equation \([16, 34]\):

\[
\frac{\partial u}{\partial t} = \frac{1}{\varepsilon} (u - u^2 - (fv + \phi) \frac{u - q}{u + q}) + D_u \nabla^2 u
\]

\[
\frac{\partial v}{\partial t} = u - v.
\]

The variable \( u \) is abstracted as a local density of plasmodium’s protoplasm and \( v \) reflects local concentration of metabolites and nutrients. We integrate the system using Euler method with five-node Laplasian operator, time step \( \Delta t = 5 \cdot 10^{-3} \) and grid point spacing \( \Delta x = 0.25 \), with the following parameters: \( \phi = \phi_0 - \eta/2 \), \( A = 0.0011109 \), \( \phi_0 = 0.0766 \), \( \varepsilon = 0.03 \), \( f = 1.4 \), \( q = 0.022 \).

Parameters \( q \) and \( f \) are inherited from model of Belousov-Zhabotinsky medium, \( \phi \) is proportional to local concentration of attractants and repellents. The parameter \( \eta \) corresponds to a gradient of chemo-attractants emitted by data planar points. The medium is perturbed by an initial excitation, where a 11 × 11 sites are assigned \( u = 1.0 \) each. The perturbation generates a wave-fragment alike propagation of protoplasm which travels along gradient \( \eta \). Repellents emitted by sites of \( P \) show very limited diffusion. Therefore repellents can be regarded as impassable obstacles (Fig. 6).

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**FIGURE 3.** Computation of concave hulls of point set $\mathbf{P}$. Images (ac) and their binarizations (bd) taken 12 h (ab) and 24 h (cd) hours after inoculation. (e) Scheme of the plasmodium interaction with $\mathbf{P}$: points/pills of $\mathbf{P}$ are circles; diffusing repellents, visible as black halo’s surrounding pills in (a) and (c), are shown by grey gradients; feeding boundary of protoplasm is marked by dotted line; major protoplasmic tube, which approximates $CH(\mathbf{P})$, is shown by solid line.
To imitate formation of the protoplasmic tubes we store values of $u$ in matrix $L$, which is processed at the end of simulation. For any site $x$ and time step $t$ if $u_x > 0.1$ and $L_x = 0$ then $L_x = 1$. The matrix $L$ represents time lapse superposition of propagating wave-fronts (Fig. 6b). The simulation is considered completed when propagating pattern envelopes $P$ and halts any further motion. At the end of simulation we repeatedly apply the erosion operation [2] (which represents a stretch-activation effect [20] necessary for formation of plasmodium tubes) to $L$. The resultant protoplasmic network (Fig. 6c) provides a good phenomenological match for networks recorded in laboratory experiments.
4. DISCUSSION

In last five years we witnessed a substantial progress in design and fabrication of working prototypes of amorphous biological computing devices based on foraging behaviour of slime mould *P. polycephalum*. Experimental laboratory prototypes of Physarum-based processors are developed for computation of minimum spanning tree, relative neighbourhood graph, Delaunay triangulation, and Voronoi diagram, see overview in [2]. By present experimental demonstration of planar shape (convex and concave hulls) computing by *P. polycephalum* we draw a logic conclusion of the years of studies into computational properties of the slime mould with regards to problems of computational geometry.
We used Kalms Tablet and Kalms Sleep pills (both brands cause similar effects to plasmodium behaviour) to represent data points. We exploited two properties of the pills. First, the pills contain Valerian and hops extracts, which are long distant attractants for *P. polycephalum* [6] and thus ‘pull’ plasmodium towards the date set. Second, the pills contain short-distance, slow diffusing components (potential candidates are magnesium stearate, stearid acid and titanium dioxide present, and more likely sucrose [35]) which may play a role of repellents preventing plasmodium from spreading inside the data set. Combination of these two properties causes *P. polycephalum* to develop a network of protoplasmic tubes with major tube approximating convex hull. Exact mechanism of short-distance repelling is unclear and will be a topic of further investigations.

Plasmodium of *P. polycephalum* approximates shape of *P* by propagating simultaneously clockwise and anti-clockwise (two branches). These branches fuse when meet one another. Thus the slime mould approximates shape of *P* in time $\frac{\Pi}{2}$, where $\Pi$ is a perimeter of the shape. The time can be reduced further by $n$ by inoculating plasmodium in $n$ loci of space at the same. Indeed the slime mould based experimental computation of planar shape can not and will not compete with existing algorithms, however will contribute towards design of future parallel embedded processors made of non-linear chemical media, and also in control and navigation of amoeboid robots.

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