Image Encryption with The Cross Diffusion of Two Chaotic Maps

Ge Jiao¹,²*, Xiaojiang Peng³, Kaiwen Duan²

¹School of Environment Protection and Safety Engineering, University of South China
Hengyang, China
²College of Computer Science and Technology, Hengyang Normal University
Hengyang, China
[E-mail: jiaoge@126.com]
³Shenzhen Institutes of Advanced Technology, Chinese Academy of Sciences, Shenzhen, China
[E-mail: xj.peng@siat.ac.cn]
*Corresponding author: Ge Jiao

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Abstract

Information security has become increasingly important with the rapid development of mobile devices and internet. An efficient encryption system is a key to this end. In this paper, we propose an image encryption method based on the cross diffusion of two chaotic maps. We use two chaotic sequences, namely the Logistic map and the Chebyshev map, for key generation which has larger security key space than single one. Moreover, we use these two sequences for further image encryption diffusion which decreases the correlation of neighboring pixels significantly. We conduct extensive experiments on several well-known images like Lena, Baboon, Koala, etc. Experimental results show that our algorithm has the characteristics of large key space, fast, robust to statistic attack, etc.

Keywords: Image encryption, chaotic map, key space, NPCR, UACI

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1. Introduction

Smartphones and tablets have accounted for over 60 percent of all smart connected consumer devices, and the global smartphone installed base will grow from four billion in 2016 to more than six billion smartphones in use by 2020, according to new analysis released by IHS Markit. With the wide application of mobile devices, the security and personality of transferred images have attracted increasing attention recently.

When an image is transmitted over Internet, it may be easily accessed, illegally used or maliciously tampered [1]. To protect the user’s image data, image encryption and data hiding in images have become two main aspects in both academia and industry. Data hiding techniques are mainly used to hide user information in images for transmitting while image encryption aims to mask the content of user images. For data hiding in images, some works [1-3] introduce reversible data hiding which is useful for recovering hidden information in the receiver. As in most data hiding researches, image encryption is the first step for masking the original image. We focus our research on image encryption in this paper.

Early methods of image encryption are mainly based on off-the-shelf data encryption technologies [4, 5], such as DES (Data Encryption Standard), AES (Advanced Encryption Standard) [6], etc. However, those methods are demonstrated to be weak in efficiency and anti-attack ability due to the intrinsic features of images [5, 7, 8]. Recently, more attention is paid to chaos-based image encryption methods [9–13]. Compared to traditional non-chaotic image encryption, the chaos-based image encryption has larger key space, is faster, and is easier to implement.

In this paper, along the chaos-based pipeline, we propose an image encryption algorithm based on the cross diffusion of two chaotic maps, namely the Logistic map and the Chebyshev map. To increase the security key space, we first run the Logistic and the Chebyshev chaotic systems starting from two random initial keys with enough iterations, and then put the result of one of them into the other with several other iterations, finally we take the results as initial conditions of chaotic maps. Pixels are encrypted by XOR operations with chaotic sequences. To decrease the correlation of pixels, we switch the used two sequences to each other according to the parity of the pixel positions. For encryption diffusion, the current encrypted pixel (in the order of RGB) is conducted another XOR operation with the previous encrypted pixel (in the order of BGR). Compared to those methods based on single chaotic map, our method has larger security key space. And our encryption method can reduce the correlation of neighbouring pixels significantly due to the cross diffusion of two maps.

We summarize our contributions as following. First, we propose to use two chaotic sequences, namely the Logistic map and the Chebyshev map, for key generation which has larger security key space without adding large extra computation. Second, we use these two sequences for further image encryption diffusion which decreases the correlation of neighboring pixels significantly.

2. Related Work

According to the visualization of encrypted images, the image encryption algorithms can be classified into three major categories: (i) position permutation based algorithm [14](ii) value transformation based algorithm and [15, 16] (iii) visual transformation based algorithm [14].
In our work we mainly focus on the value transformation based algorithm. Along this line, there are two main categories: traditional image encryption methods and chaos-based methods.

The characteristics of low correlation, large key space, high key sensitivity, high entropy, and low time complexity are the main issues for a good image encryption algorithm. We refer to [15] for a detail survey on all these points.

Traditional image encryption methods are mainly based on standard data encryption algorithms [5, 7, 8]. Since the chaotic encryption algorithm proposed in [18], many works based on chaotic systems have been proposed [8, 19–30]. Pak and Huang proposed a simple and effective image encryption method based on the difference of outputs from two identified 1D chaotic systems [26]. Compared to single chaotic map, this method improves the chaotic performance and is more robust to any attacks. Rostami et al. proposed a parallel image encryption with chaotic windows based on logistic map [25], which has large key space and uniform histogram. This method first divides an image to image patches, and then performs XOR operation on patches with chaotic windows. The NPCR(number of pixels change rate) and UACI(unified average changing intensity) measurements show that it is effective against differential attacks. Patidar et al. presented a modified substitution diffusion algorithm using chaotic standard and logistic maps [17]. This method is very fast in permutation and diffusion, and is effective to resist plaintext attacks. Praveenkumar et al. combined chaotic and Chebyshev keys and proposed a multi-level encryption system based on the frequency domain and the chaotic pixel permutation [21]. Yoon et al. first generated a random sequence using chaotic systems, and then permuted the rows and columns of an image according to the random sequence [22]. All those methods need large computation cost in value transformation or generating random sequences, and degrade on mobile devices in encryption speed.

3. Preliminaries on chaotic systems and chaos based image encryption

Chaos-based algorithms have shown some remarkable properties in security, complexity, performance, speed etc. The following are the characteristics of chaotic maps:

1. They are deterministic, i.e. they have several mathematical equations or formulation which ruling their behavior.
2. They are sensitive to initial conditions.
3. They are unpredictable and non-linear that is a small change can produce large or huge effects.
4. They appear to be random and disorderly.
5. They usually produce fractal patterns.

Fig. 1 shows the first part of our cross diffusion of Logistic and Chebyshev. The outputs, \( x_0 \) and \( x_e \), are the initial conditions of Logistic map and Chebyshev map for the three channels of an image (i.e. RGB). Note that the second part is embeded in the image encryption algorithm where the Logistic sequence and the Chebyshev sequence are used alternately.
In fact, chaotic systems and cryptographic algorithms have many similar properties [27], such as both are sensitive to initial parameters or keys, which makes chaos-based image encryption more and more popular. There are two widely-used chaotic maps for encryption, namely the Logistic map and the Chebyshev map. The Logistic map is a classic non-linear chaotic function which has complex chaotic behavior. The formulation of Logistic map is as follows,

\[ x_{n+1} = \mu x_n (1 - x_n), \quad x_n \in (0, 1) \]  

where \( x_n \) is the value after \( n \) iterations, and \( \mu \) is the system parameter which makes the system chaotic and ensures \( x_n \) \( \in (0, 1) \) when \( \mu \in (3.5699456, 4] \).

The ChebyShev map is very sensitive to its initial conditions and can generate long-term unpredictable chaotic sequences. With increasing iterations, the generated chaotic trajectory becomes uniform and mixing, and the initial adjacent points will be separated exponentially. The formulation of ChebyShev map is as follows,

\[ x_{n+1} = \cos(k \arccos(x_n)), \quad x_n \in [-1, 1] \]  

where \( k \) is the control parameter which makes the function chaotic when \( k \geq 2 \).

**Chaos-based image encryption.** The most popular pipeline of chaos-based image encryption mainly includes three steps: (i) generating a key as initial conditions for chaotic systems, (ii) generating a chaotic sequence, and (iii) conducting XOR operations on pixels with the chaotic sequence.

**4. The proposed image encryption approach**

We propose an image encryption approach based on the cross diffusion of the Logistic map and the ChebyShev map. We present the cross diffusion scheme and the image encryption algorithm in the following two subsections.

**4.1 The cross diffusion of Logistic and ChebyShev**

The cross diffusion scheme is shown in two parts. Figure 1 illustrates the first part where the cross diffusion scheme is mainly used for generating initial conditions. Since we apply two chaotic maps for color image encryption, we need initial conditions for both maps of three channels (i.e. RGB).
Fig. 2 shows our image encryption algorithm.

\begin{algorithm}
\caption{Our image encryption algorithm}
\begin{algorithmic}[1]
\Require $I \in \mathbb{R}^{M \times N}$: the original image; $X_0^{cl}$: the initial condition of Chebyshev map for RGB channels; $X_0^{lc}$: the initial condition of Logistic map for RGB channels;
\Ensure the encrypted image $I^*$
\For{each pixel $P_i \in \mathbb{R}^3$, $i \in [0, M \times N]$}
\If{$i$ is even}
\State generate $X_{i+1}^{cl}$ using Chebyshev function;
\State $P_i^{*} = X_{i+1}^{cl} \oplus P_i$;
\If{$i > 0$}
\State $P_i^{*} = P_i^{*} \oplus P_{i-1}^{*}$, where $P_i^{*}$ is the ciphertext of $P_i$;
\Else
\State $P_i^{*} = P_i^{*}$;
\EndIf
\Else
\State generate $X_{i+1}^{lc}$ using Logistic function;
\State $P_i^{*} = X_{i+1}^{lc} \oplus P_i$;
\EndIf
\EndFor
\end{algorithmic}
\end{algorithm}

Fig. 2. Our image encryption algorithm

For the Chebyshev map, starting from a Key $c \in [1, 1]$, we first conduct N iterations (we fix N to 100 in our implementation) using the Chebyshev function, and then input the absolute value of $X_i$ into the Logistic function, finally we generate three initial conditions by another 3 iterations. For the Logistic map, starting from a Key $l \in (0, 1)$, we first perform N iterations using the Logistic function, and then input the result to the Chebyshev function, finally we also generate three initial conditions by another 3 iterations.

4.2 The image encryption algorithm

Our image encryption method is described in Algorithm 1. The inputs of our image encryption methods are the outputs from Figure 1 and an original image I. Given a pixel $P_i \in I$ with position $i$, the main steps of our approach is as follows:

1. Check the parity of $i$, go to (2) if it is even, otherwise go to (4).

2. Compute $\chi_l^{ci}$, by the Chebyshev function (i.e. equation (2)), conduct XOR operation with $\chi_l^{ct}$ and $P_i$ for diffusion, and denote the result as $P_i^*$. 

3. Perform another XOR operation with $P_i^*$ and the previous encrypted pixel $P_{i-1}^*$ if $i > 0$, since $P_{i-1}^*$ is obtained by the other chaotic map.

4. Compute $\chi_l^{te}$, by the Logistic function (i.e. equation (1)), conduct XOR operation with $\chi_l^{tc}$ and $P_i$ for diffusion, and denote the result as $P_i^*$. 

5. Perform another XOR operation with $P_i^*$ and the previous encrypted pixel $P_{i+1}^*$, and get the final target pixel $P_i^*$.
Fig. 3 shows several examples using our image encryption algorithm. The encrypted images are noise-like images as expected. We use these 8 images (a-h) for our experiments and present extensive analysis in the next section.
To evaluate our image encryption algorithm, we conduct extensive experimental analysis on several standard images with varied sizes. In this section, we make time analysis, statistical analysis, key space analysis, information entropy analysis, and sensitivity analysis.
5.1 Time analysis

Regardless the security, the time delay is the most important user experience for mobile devices. For the time analysis, we conduct experiments on the Lena.bmp image with varied scales, and other standard images such as Girl.bmp, Baboon.bmp, and Barbara.bmp. Table 1 shows a comparison between our algorithm and other algorithm designed for mobile platform [23, 28]. We compute the average running time after 100 runs using Opencv2.4.13 (Android version). The running time is obtained on a smartphone with Android6.0 OS, 1.7 GHz Hisilicon Kirin920 CPU, and 3 GB RAM. From Table 1, we see that our algorithm is very efficient in both encryption and decryption.

| Image size | [23](s) | [28](s) | Ours(s) |
|------------|---------|---------|---------|
| Lena.bmp (128×128) | 0.070 | 0.068 | - | - | 0.0039 | 0.0033 |
| Lena.bmp 256×256) | 0.261 | 0.269 | - | - | 0.0152 | 0.0130 |
| Lena.bmp (512×512) | 0.298 | 0.287 | - | - | 0.0621 | 0.0524 |
| Lena.bmp (1366×768) | 1.175 | 1.127 | - | - | 0.2460 | 0.2007 |
| Girl.bmp (256×256) | 0.198 | 0.201 | 0.047 | 0.050 | 0.0355 | 0.0085 |
| Baboon.bmp (512×512) | 0.256 | 0.249 | 0.181 | 0.180 | 0.0868 | 0.0347 |
| Barbara.bmp (787×576) | 0.446 | 0.449 | 0.306 | 0.312 | 0.0992 | 0.0658 |

5.2 Statistical analysis

We mainly analyze the histograms and the correlation of two adjacent pixels using the Baboon.bmp image in this section.

5.2.1 Histogram analysis

Image histogram describes the distribution of pixel values of an image. Flatter is better to resist statistic attacks for the image histogram. We conduct the histogram analysis on the Baboon.bmp image. Figure 4 shows the histograms of the original image and the encrypted image. The histogram comparison of RGB channels (from top to down) on Lena and Baboon. In each subfigures, the original image histograms are shown in the left, and the encrypted image histograms in the right. From Figure 4, we see that all the channels of the encrypted image have good uniform distributions, thus it is effective to resist statistic attacks.
5.2.2 Correlation analysis of two adjacent pixels

Neighbouring pixels in an image are strongly correlated in the directions of vertical, horizontal and diagonal. This correlation of encrypted images can be used by statistical attacks. To this end, an encryption algorithm should be able to remove the correlation of adjacent pixels.

The correlation coefficient is calculated as follows:

\[ r_{xy} = \frac{\text{cov}(x, y)}{\sqrt{D(x)} \sqrt{D(y)}} \]

where \( \text{cov}(x, y) = \frac{1}{N} \sum_{i=1}^{N} (x_i - E(x))(y_i - E(y)) \)

\[ D(x) = \frac{1}{N} \sum_{i=1}^{N} (x_i - E(x))^2, \quad E(x) = \frac{1}{N} \sum_{i=1}^{N} x_i \]

where \( x \) and \( y \) are the color values of two neighbouring pixels in images, \( N \) is the number of selected pixel pairs, and \( r_{xy} \) is in \([-1, 1]\). It is strongly related if \(|r_{xy}| > 0.8\), otherwise weakly related if \(|r_{xy}| < 0.3\).

We conduct correlation analysis on three standard images, namely Cameraman.bmp, Lena.bmp, and flower.bmp. The correlation coefficients among adjacent pixels at horizontal, vertical and diagonal directions on these images are shown in Table 2. Those results are obtained by randomly selecting 3000 pixel pairs in corresponding directions. From Table 2, we see that adjacent pixels are strongly correlated in the original image and almost uncorrelated in the encrypted image.

![Figure 4. Comparison of histogram of original image and encrypted image](image-url)
encrypted image. The last two rows in Table 2 compare our method with a recent chaos based method. Our method shows smaller correlation coefficients in average.

Table 2. The comparison of correlation coefficient
The results on Lena.bmp and flower.bmp are not available in [25]

| Method  | Image            | Channel | Original image | Encrypted image |          |          |          |
|---------|------------------|---------|---------------|----------------|--------|--------|--------|
|         |                  |         | Vertical | Horizontal | Diagonal | Vertical | Horizontal | Diagonal |          |        |        |        |
| Ours    | Cameraman.bmp p(256×256) | Red     | 0.96148 | 0.93074 | 0.91664 | 0.01722 | 0.02364 | 0.01949 |          |        |        |        |
|         |                  | Green   | 0.96148 | 0.93074 | 0.91664 | 0.03334 | 0.01521 | 0.03959 |          |        |        |        |
|         |                  | Blue    | 0.96148 | 0.93074 | 0.91664 | 0.00992 | 0.02257 | 0.00792 |          |        |        |        |
|         |                  | Average | -       | -        | -        | 0.02016 | 0.02047 | 0.02233 |          |        |        |        |
| [25]    | Cameraman.bmp p    | Average | -       | -        | -        | 0.02608 | 0.02424 | 0.02446 |          |        |        |        |
| Ours    | Lena.bmp(256×256)  | Red     | 0.96860 | 0.94530 | 0.92132 | 0.01425 | 0.02421 | 0.02745 |          |        |        |        |
|         |                  | Green   | 0.97094 | 0.94633 | 0.92572 | 0.00557 | 0.00061 | 0.02194 |          |        |        |        |
|         |                  | Blue    | 0.94140 | 0.90402 | 0.87477 | 0.01317 | 0.00166 | 0.00498 |          |        |        |        |
|         |                  | Average | -       | -        | -        | 0.01099 | 0.00883 | 0.01812 |          |        |        |        |
| Ours    | Flower.bmp(120×761) | Red     | 0.99452 | 0.99467 | 0.99042 | 0.03383 | 0.01811 | 0.01331 |          |        |        |        |
|         |                  | Green   | 0.99454 | 0.99493 | 0.99106 | 0.01992 | 0.01440 | 0.00373 |          |        |        |        |
|         |                  | Blue    | 0.98617 | 0.98836 | 0.97799 | 0.00774 | 0.00118 | 0.00737 |          |        |        |        |
|         |                  | Average | -       | -        | -        | 0.02050 | 0.01123 | 0.00814 |          |        |        |        |

5.3 Key space analysis
The key space is a basic measurement of any encryption algorithms. In our algorithm we have 8 keys, namely $Key_c$, $Key_l$, $x_0c$, $x_1c$, $x_2c$, $x_0l$, $x_1l$, and $x_2l$. $Key_c$ and $Key_l$ are computed in the accuracy of $10^{-16}$. $x_0c$, $x_1c$, $x_2c$ are generated simultaneously, and the joint key space of them is $3 \times 10^{16}$. The key space of [$x_0l$, $x_1l$, $x_2l$] is $3 \times 10^{16}$ as well. So the total key space is $10^{16} \times 3 \times 10^{16} \times 3 \times 10^{16} \approx 2^{218}$. It means our algorithm has large enough key space to withstand the brute force attack. And our key space is significantly larger than that ($2^{138}$) in a recent chaos based algorithm [26].

5.4 Information entropy analysis
Information entropy refers to disorder or uncertainty [29]. We compute the image information as follows,

$$H = - \sum_{i=0}^{255} p_i \log p_i$$

(4)

where $p_i$ is the probability of gray value $i$ in an image, which can be obtained from the histogram. Table 3 shows the image information entropies of both the original image (Baboon.bmp) and the encrypted image. A uniform distribution has the maximum value 8 in
Equation (4). Our encrypted image has larger entropy than the original one, which means the encrypted image is more random.

**Table 3.** The information entropy analysis

| Image        | Original image | Encrypted image |
|--------------|---------------|-----------------|
|              | R  | G  | B  | R  | G  | B  |
| Baboon(512*512) | 7.64129 | 7.34797 | 7.66837 | 7.99747 | 7.99748 | 7.99721 |
| Lena(512*512)   | 7.25310 | 7.59404 | 6.96843 | 7.99929 | 7.99932 | 7.99922 |
| Flower(1200*761)| 7.43894 | 7.64537 | 7.12677 | 7.99976 | 7.99979 | 7.99978 |

5.5 Sensitivity analysis

We make qualitative and quantitative sensitivity analysis in our experiment. For qualitative analysis, we add a small noise ($10^{-16}$) to both $Key_c$ (we use 0.456) and $Key_l$ (we use 0.456), and show the decrypted images in Fig. 5. The very different decrypted images indicate the good sensitivity of our algorithm.

The sensitivity also can be quantitatively evaluated by NPCR(number of pixels change rate) and UACI(unified average changing intensity). We compute those two values from two encrypted images $c_1$ (with $Key_c=Key_l=0.456$) and $c_2$ (with $Key_c=Key_l=0.456+10^{-16}$). NPCR is formulated as follows,

$$NPCR = \frac{1}{M \times N} \sum_{i=1}^{M} \sum_{j=1}^{N} D(i, j) \times 100\%$$

where $D(i, j) = 0$ if $c_1(i, j) = c_2(i, j)$, otherwise $D(i, j) = 1$. And the UACI is formulated as follows,

$$UACI = \frac{1}{M \times N} \sum_{i=1}^{M} \sum_{j=1}^{N} \left| \frac{c_1(i, j) - c_2(i, j)}{255} \right| \times 100\%.$$  

Table 4 shows NPCR and UACI measurements on four images. In general, larger values indicate that the algorithm is more sensitive and better. We compare our algorithm to the latest chaos-based method [26] with similar initial keys and the same small noise (i.e. $10^{-16}$).
From the comparison, our algorithm presents better performance in NPCR and UACI measurements.

Table 4. The NPCR and UACI in three channels. Larger is better

| Image   | Method | NPCR(%) | UACI(%) |
|---------|--------|---------|---------|
|         |        | R       | G       | B       | R       | G       | B       |
| Lena.bmp | [26]   | 99.6552 | 99.6277 | 99.588  | 33.4846 | 33.4132 | 33.3441 |
|         | ours   | 99.6561 | 99.6309 | 99.6063 | 33.5069 | 33.4212 | 33.5187 |
| Koala.bmp | [26] | 99.5834 | 99.6262 | 99.5926 | 33.5202 | 33.4907 | 33.4115 |
|         | ours   | 99.6249 | 99.6365 | 99.6173 | 33.5293 | 33.4928 | 33.4234 |
| Flower.bmp | [26] | 99.6277 | 99.6063 | 99.5575 | 33.3951 | 33.5272 | 33.4008 |
|         | ours   | 99.6315 | 99.6127 | 99.6276 | 33.4502 | 33.3814 | 33.5937 |
| Dog.bmp | [26]   | 99.6353 | 99.6063 | 99.5926 | 33.4502 | 33.3814 | 33.5937 |
|         | ours   | 99.6396 | 99.6127 | 99.6276 | 33.4502 | 33.3814 | 33.5937 |

To evaluate the sensitivity, we also mimic differential attacks as in [25,31]. The image was primarily encrypted. Next, we changed one pixel in the plain image and encrypted it. Then the NPCR and UACI between the corresponding cipher images are calculated. This process was iterated 100 times by changing the random pixels. We conduct extensive experiments on the widely-used USC-SIPI image processing dataset. A comprehensive comparison of NPCR and UACI are shown in Table 5. As shown in the last row of Table 5, our method is superior or comparable to [25,31] in both mean NPCR and UACI.

Table 5. NPCR and UACI values of encrypted images for the Proposed Method and other algorithms running a hundred times on USC-SIPI image dataset

| Image   | Method | NPCR(%) | UACI(%) |
|---------|--------|---------|---------|
|         |        | Max     | Min     | Mean   | Max     | Min     | Mean   |
| 5.1.09  | [25]   | 99.68262 | 99.54376 | 99.61049 | 33.69240 | 33.20391 | 33.42682 |
|         | [31]   | 99.66736 | 99.53461 | 99.60907 | 33.73600 | 33.29381 | 33.48055 |
|         | ours   | 99.67874 | 99.54415 | 99.60912 | 33.67090 | 33.21256 | 33.44120 |
| 5.1.10  | [25]   | 99.65515 | 99.55597 | 99.61079 | 33.71333 | 33.29039 | 33.47612 |
|         | [31]   | 99.69330 | 99.54681 | 99.61218 | 33.69135 | 33.18796 | 33.44003 |
|         | ours   | 99.67041 | 99.55798 | 99.61112 | 33.72295 | 33.21158 | 33.46348 |
| 5.1.11  | [25]   | 99.65668 | 99.53461 | 99.60550 | 33.67225 | 33.25851 | 33.47719 |
|         | [31]   | 99.65820 | 99.53461 | 99.60997 | 33.71989 | 33.21694 | 33.46591 |
|         | ours   | 99.67041 | 99.53579 | 99.60962 | 33.69273 | 33.21263 | 33.47443 |
| 5.1.12  | [25]   | 99.66125 | 99.55444 | 99.60843 | 33.62589 | 33.25152 | 33.45952 |
|         | [31]   | 99.65820 | 99.54681 | 99.60915 | 33.64215 | 33.19594 | 33.43716 |
|         | ours   | 99.67521 | 99.54734 | 99.60912 | 33.70610 | 33.22198 | 33.47070 |
| 5.1.13  | [25]   | 99.67041 | 99.53918 | 99.60722 | 33.67809 | 33.20683 | 33.45420 |
|         | [31]   | 99.66888 | 99.55750 | 99.61072 | 33.66077 | 33.28392 | 33.47259 |
|         | ours   | 99.67041 | 99.54028 | 99.60912 | 33.76906 | 33.20987 | 33.48710 |
| 5.2.08  | [25]   | 99.64409 | 99.58458 | 99.61077 | -        | -        | -        |
|         | [31]   | 99.63646 | 99.58496 | 99.60844 | -        | -        | -        |
|         | ours   | 99.64218 | 99.57390 | 99.60872 | 33.60928 | 33.38309 | 33.48207 |
| 5.2.09  | [25]   | 99.63455 | 99.57886 | 99.61013 | -        | -        | -        |
|         | [31]   | 99.63951 | 99.57695 | 99.60905 | -        | -        | -        |
|         | ours   | 99.64218 | 99.57390 | 99.60872 | 33.60068 | 33.34674 | 33.46001 |
5.2.10

|      | [25] | [31] | ours |
|------|------|------|------|
|      | 99.64943 | 99.58649 | 99.60957 |
|      | 99.63837 | 99.57504 | 99.60878 |
| ours | 99.64218 | 99.57390 | 99.60872 |
|      | 33.57785 | 33.34855 | 33.46840 |

5.3.01

|      | [25] | [31] | ours |
|------|------|------|------|
|      | 99.62454 | 99.59812 | 99.60938 |
|      | 99.62435 | 99.59164 | 99.60940 |
| ours | 99.62330 | 99.59421 | 99.60950 |
|      | 33.51622 | 33.39734 | 33.45913 |

5.3.02

|      | [25] | [31] | ours |
|------|------|------|------|
|      | 99.63417 | 99.59249 | 99.60966 |
|      | 99.62502 | 99.59478 | 99.60947 |
| ours | 99.62330 | 99.59421 | 99.60950 |
|      | 33.52143 | 33.39220 | 33.46750 |

7.1.01

|      | [25] | [31] | ours |
|------|------|------|------|
|      | 99.64523 | 99.57809 | 99.60949 |
|      | 99.63531 | 99.58115 | 99.60944 |
| ours | 99.64218 | 99.57390 | 99.60872 |
|      | 33.59882 | 33.35273 | 33.47480 |

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|      | [25] | [31] | ours |
|------|------|------|------|
|      | 99.63646 | 99.58305 | 99.60839 |
|      | 99.64027 | 99.56818 | 99.60872 |
| ours | 99.64218 | 99.57390 | 99.60947 |
|      | 33.57845 | 33.34303 | 33.47754 |

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|      | [25] | [31] | ours |
|------|------|------|------|
|      | 99.63303 | 99.58191 | 99.60969 |
|      | 99.64180 | 99.57008 | 99.60872 |
| ours | 99.64218 | 99.57390 | 99.60872 |
|      | 33.59614 | 33.36760 | 33.47169 |

gray21_512

|      | [25] | [31] | ours |
|------|------|------|------|
|      | 99.63951 | 99.57504 | 99.60947 |
|      | 99.63875 | 99.58191 | 99.60103 |
| ours | 99.64218 | 99.57465 | 99.61012 |
|      | 33.59614 | 33.36760 | 33.47169 |

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|      | [25] | [31] | ours |
|------|------|------|------|
|      | 99.64600 | 99.58458 | 99.60777 |
|      | 99.63303 | 99.58305 | 99.60969 |
| ours | 99.64218 | 99.57390 | 99.60872 |
|      | 33.53510 | 33.37913 | 33.47222 |

testpat. 1k

|      | [25] | [31] | ours |
|------|------|------|------|
|      | 99.62378 | 99.59635 | 99.60941 |
|      | 99.62206 | 99.59440 | 99.60900 |
| ours | 99.62329 | 99.59421 | 99.60949 |
|      | 33.53071 | 33.41411 | 33.46971 |

Camer aman

|      | [25] | [31] | ours |
|------|------|------|------|
|      | 99.66583 | 99.55750 | 99.61101 |
|      | 99.68109 | 99.55292 | 99.60669 |
| ours | 99.66857 | 99.54087 | 99.61075 |
|      | 33.71245 | 33.29961 | 33.50756 |

total results

|      | [25] | [31] | ours |
|------|------|------|------|
|      | 99.68262 | 99.53461 | 99.60917 |
|      | 99.69330 | 99.53461 | 99.60966 |
| ours | 99.67874 | 99.53579 | 99.60935 |
|      | 33.76906 | 33.20987 | 33.46754 |

5. Conclusion

This paper proposed a new image encryption algorithm based on the cross diffusion of two chaotic maps. The proposed algorithm has four main properties. First, it has very large security key space due to the cross diffusion of Logistic map and ChebyShev map in key generation. Second, it makes the correlation of adjacent pixels very weak due to the cross diffusion of Logistic map and ChebyShev map in encryption process. Third, it is very sensitive to tiny differences of the initial condition due to the XOR operation between previous and current pixels. Forth, it is realtime on mobile devices.

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Ge Jiao received his MA.Eng from Hunan University, Changsha, China, in 2010. He is currently working toward the PhD degree in safety engineering at University of South China. He works at Hengyang Normal University as an associate professor. His research interest focuses on information security, image encryption and encryption chip side channel attacks and protection.

Xiaojiang Peng received his master degree in school of Information Science and Technology from Southwest Jiaotong University in 2014. He currently is an Associate Professor at the Multimedia Lab, Shenzhen Institutes of Advanced Technology, Chinese Academy of Sciences, China. He was a postdoctoral researcher at Idiap Institute, Switzerland from 2016 to 2017, and was a postdoctoral researcher in LEAR Team, INRIA, France, working with Prof. Cordelia Schmid from 03/2015 to 07/2016. He serves as a reviewer for IEEE Transactions on Image Processing, IEEE Transactions on Multimedia, Image and Vision Computing, Machine Vision and Applications, IEEE Signal Processing Letter, Multimedia Tools and Applications, Neurocomputing, IET Computer Vision, FG, etc. His research focus is in the area of image processing, action recognition and detection, face recognition and deep learning.

Kaiwen Duan is currently pursuing a bachelor degree at the College of Information Engineering in Hengyang Normal University, Hengyang, China. His current research interests include the image processing and information security.