Generalized entanglement monogamy and polygamy relations for $N$-qubit systems

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We investigate the generalized monogamy and polygamy relations $N$-qubit systems. We give a general upper bound of the $\alpha$th ($0 \leq \alpha \leq 2$) power of concurrence for $N$-qubit states. The monogamy relations satisfied by the $\alpha$th ($0 \leq \alpha \leq 2$) power of concurrence are presented for $N$-qubit pure states under the partition $AB$ and $C_1...C_{N-2}$, as well as under the partition $ABC_1$ and $C_2...C_{N-2}$. These inequalities give rise to the restrictions on entanglement distribution and the trade off of entanglement among the subsystems. Similar results are also derived for negativity.

PACS numbers:

INTRODUCTION

Quantum entanglement, as a physical resource in quantum communication and quantum information processing, has been the subject of many studies in recent years. The study of quantum entanglement from different points of view has been a very active area and has led to many interesting results. Characterizing and quantifying entanglement have emerged as one of the prominent themes of quantum information theory. Although bipartite entanglement is relatively well understood at least for qubit systems, the characterization or classification of multipartite entanglement is still very challenging.

Monogamy is one of the key features of multipartite entanglement. As one of the fundamental differences between quantum entanglement and classical correlations, a key property of entanglement is that a quantum system entangled with one of other systems limits its entanglement with the remaining ones. The monogamy relations give rise to the restrictions on distribution of entanglement in the multipartite setting. It is also the key ingredient behind secure quantum cryptography and has played an important role in condensed matter physics such as the $n$-representability problem for fermions.

Thus, it is an important task to understand entanglement monogamy and its characterization in order to unveil the power of quantum entanglement in multipartite systems. Monogamy of entanglement was first noted by Coffman, Kundu, and Wootters in terms of an inequality satisfied by the squared concurrence $C$, often referred to as the CKW inequality. Concurrence is a well-known measure of bipartite entanglement, introduced by Wootters.

For a tripartite system $A$, $B$ and $C$, the usual monogamy of an entanglement measure $\mathcal{E}$ implies that the entanglement between $A$ and $BC$ satisfies $\mathcal{E}_{ABC} \geq \mathcal{E}_{AB} + \mathcal{E}_{AC}$. For a pure three-qubit state $|\psi\rangle_{ABC}$, the CKW inequality is first given by $C^2(\rho_{ABC}) \geq C^2(\rho_{AB}) + C^2(\rho_{AC})$, where $A, B, C$ are three qubits with $A$ as the focus qubit, and $\rho_{AB}$ stands for the bipartite splitting between $A$ and $BC$. The concurrence of a two-qubit state $\rho$ is given by $C(\rho) = \max \{0, \lambda_1 - \lambda_2 - \lambda_3 - \lambda_4\}$, where $\lambda_1, \lambda_2, \lambda_3, \lambda_4$ are the square root of the eigenvalues of the matrix $\rho(\sigma_y \otimes \sigma_y)\rho^*(\sigma_y \otimes \sigma_y)$ in decreasing order, $\sigma_y$ is the Pauli matrix, and $\rho^*$ denotes the complex conjugate of $\rho$.

IMPROVED GENERALIZED MONOGAMY AND POLYGAMY RELATIONS OF CONCURRENCE

Let $H_X$ denote the finite dimensional vector space associated with qubit $X$. For a bipartite pure state $|\psi\rangle_{AB}$ in vector space $H_A \otimes H_B$, the concurrence is given by

$$C(|\psi\rangle_{AB}) = \sqrt{2[1 - \text{Tr}(\rho_A^2)]},$$

(1)
where \( \rho_A \) is the reduced density matrix by tracing over the subsystem \( B, \rho_A = \text{Tr}_B(\psi_{AB} \langle \psi |) \). The concurrence for a bipartite mixed state \( \rho_{AB} \) is defined by the convex roof
\[
C(\rho_{AB}) = \min_{\{p_i, |\psi_i\rangle\}} \sum_i p_i C(|\psi_i\rangle),
\]
where the minimum is taken over all possible decompositions of \( \rho_{AB} = \sum p_i |\psi_i\rangle \langle \psi_i| \), with \( p_i \geq 0 \), \( \sum_i p_i = 1 \) and \( |\psi_i\rangle \in H_A \otimes H_B \).

For a tripartite state \( |\psi\rangle_{ABC} \), the concurrence of assistance (COA) is defined by [19]
\[
C_a(|\psi\rangle_{ABC}) = C_a(\rho_{AB}) = \max_{\{p_i, |\psi_i\rangle\}} \sum_i p_i C(|\psi_i\rangle),
\]
for all possible ensemble realizations of \( \rho_{AB} = \text{Tr}_C(\rho_{ABC} |\psi\rangle \langle \psi|) = \sum_i p_i |\psi_i\rangle \langle \psi_i| \). When \( \rho_{AB} \) is a pure state, one has
\[C(|\psi\rangle_{ABC}) = C_a(\rho_{AB}).\]

For an \( N \)-qubit state \( |\psi\rangle_{AB_1 \ldots B_{N-1}} \in H_A \otimes H_{B_1} \otimes \cdots \otimes H_{B_{N-1}} \), the concurrence \( C(|\psi\rangle_{AB_1 \ldots B_{N-1}}) \) of the state \( |\psi\rangle_{AB_1 \ldots B_{N-1}} \), viewed as a bipartite with partitions \( A \) and \( B_1 B_2 \cdots B_{N-1} \), satisfies the monogamy inequality for \( \alpha \geq 2 \) [20],
\[
C^\alpha(|\psi\rangle_{AB_1 B_2 \cdots B_{N-1}}) \geq C^\alpha(\rho_{AB}) + C^\alpha(\rho_{AB_2}) + \cdots + C^\alpha(\rho_{AB_{N-1}}), \tag{2}
\]
where \( C(\rho_{AB_i}) \) is the concurrence of \( \rho_{AB_i} = \text{Tr}_{B_{1 \ldots i-1} B_{i+1 \ldots N-1}}(\rho) \).

The dual inequality satisfied by COA for \( N \)-qubit states has the form [21],
\[
C^\alpha(|\psi\rangle_{A|B_1 B_2 \cdots B_{N-1}}) \leq \sum_{i=1}^{N-1} C^\alpha(\rho_{AB_i}). \tag{3}
\]

Furthermore, the authors in [22] give a generalized monogamy relation for \( \alpha \geq 2 \),
\[
C^\alpha(\rho_{AB_1 B_2 \cdots B_{N-1}}) \geq C^\alpha(\rho_{AB_1}) + C^\alpha(\rho_{AB_2}) + \cdots + C^\alpha(\rho_{AB_{N-1}}). \tag{4}
\]
However, there are no dual inequalities yet satisfied by the athen power of COA for \( N \)-qubit states. In this paper, we will give some monogamy and polygamy relations for \( N \)-qubit states in terms of the athen power of COA.

The concurrence \( (1) \) is related to the linear entropy \( T(\rho) \) of a state \( \rho \), \( T(\rho) = 1 - \text{Tr}(\rho^2) \) [22]. For a bipartite state \( \rho_{AB} \), \( T(\rho) \) has the property [23]
\[
|T(\rho_A) - T(\rho_B)| \leq T(\rho_{AB}) \leq T(\rho_A) + T(\rho_B). \tag{4}
\]

For convenience, we write (5) as follows
\[
C^\alpha(|\psi\rangle_{AB_1 B_2 \cdots B_{N-1}}) \leq \sum_{i=1}^k C^\alpha(\rho_{A M_i}), \tag{5}
\]
where \( C^\alpha(\rho_{A M_i}) = \sum_{j=M_i+1}^{M_0} C^2(\rho_{A B_i}) \) with \( M_0 = 0, \sum_{i=1}^k M_i = N - 1 \). The summation on the right hand side of (5) has been separated into \( k \) parts. There is always a choice of \( M_i \), such that the above relations is true.

**Theorem 1.** For any \( N \)-qubit pure state \( |\psi\rangle_{A|B_1 B_2 \cdots B_{N-1}} \), we have
\[
C^\alpha(|\psi\rangle_{A|B_1 B_2 \cdots B_{N-1}}) \leq C^\alpha(\rho_{A M_i}) + \frac{\alpha}{2} C^\alpha(\rho_{A M_2}) + \cdots + \frac{(\alpha/2)^{k-1}}{C^\alpha(\rho_{A M_k})}, \tag{6}
\]
for all \( 0 \leq \alpha \leq 2 \).

**Proof.** Without loss of generality, we can always assume that \( C^2(\rho_{A M_i}) \geq \sum_{i=t+1}^{k} C^2(\rho_{A M_i}), 1 \leq t \leq k - 1, 2 \leq k \leq N - 1 \), by reordering \( M_1, M_2, \cdots, M_k \) and/or relabeling the subsystems in need. Form the result in [21], we have
\[
C^\alpha(|\psi\rangle_{A|B_1 B_2 \cdots B_{N-1}}) \leq \left(C^\alpha(\rho_{A M_1}) + \sum_{i=2}^{k} C^\alpha(\rho_{A M_i})\right)^{\frac{\alpha}{2}} \leq \frac{\alpha}{2} + \frac{(\alpha/2)^{k-1}}{C^\alpha(\rho_{A M_k})}. \tag{7}
\]
where the first inequality is due to [5]. Using the inequality \( (1 + t)^{\frac{\alpha}{2}} \leq 1 + xt \leq 1 + xt^{\alpha} \) for \( 0 \leq x \leq 1, 0 \leq t \leq 1 \), we get the second inequality. \( \square \)

Theorem 1 gives the polygamy relation of the athen power of concurrence for \( N \)-qubit pure state \( |\psi\rangle_{A|B_1 B_2 \cdots B_{N-1}} \), based on the COA. For the case of \( k = N - 1 \), we have the following result,
\[
C^\alpha(|\psi\rangle_{A|B_1 B_2 \cdots B_{N-1}}) \leq C^\alpha(\rho_{A B_1}) + \frac{\alpha}{2} C^\alpha(\rho_{A B_2}) + \cdots + \frac{(\alpha/2)^{N-2}}{C^\alpha(\rho_{A B_{N-1}})}. \tag{8}
\]
Specially, for \( \alpha = 2 \), (6) or (8) reduces to the result in [21].

**Example 1.** Let us consider the three-qubit state \( |\psi\rangle \), which can be written in the generalized Schmidt decomposition form [25, 29],
\[
|\psi\rangle = \lambda_0 |000\rangle + \lambda_1 e^{i\varphi} |100\rangle + \lambda_2 |101\rangle + \lambda_3 |110\rangle + \lambda_4 |111\rangle,
\]

where \( \lambda_0, \lambda_1, \lambda_2, \lambda_3, \lambda_4 \) are real numbers. Without loss of generality, we can always assume that \( \lambda_0^2 \geq \sum_{i=1}^{k} \lambda_i^2 \). Then we have
\[
C^\alpha(|\psi\rangle) \leq \lambda_0^\alpha + \frac{\alpha}{2} \lambda_1^\alpha + \cdots + \frac{(\alpha/2)^{k-1}}{\lambda_0^\alpha}. \tag{9}
\]
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\]
where \( \lambda_i \geq 0, \ i = 0, \ldots, 4 \) and \( \sum_{i=0}^4 \lambda_i^2 = 1 \). We have \( C(\rho_{ABC}) = 2\lambda_0\sqrt{\lambda_2^2 + \lambda_3^2 + \lambda_4^2} \), \( C(\rho_{AB}) = 2\lambda_0\lambda_2 \), \( C(\rho_{AC}) = 2\lambda_0\lambda_3 \), \( C(\rho_{BC}) = 2\lambda_0\lambda_3 \), \( C(\rho_{AB}) = 2\lambda_0\lambda_2 \), and \( C(\rho_{AC}) = 2\lambda_0\lambda_3 \). Set \( \lambda_0 = \lambda_1 = \lambda_3 = \lambda_4 = \sqrt{\frac{2}{3}} \). One gets \( C(\rho_{ABC}) = (\frac{2\sqrt{2}}{3})^\alpha, \ C(\rho_{AB}) = (1 + \frac{\alpha}{2}) (\frac{2\sqrt{2}}{3})^\alpha \). The relation from (6) is shown in Fig. 1.

[Remark 1.] We have obtained a general polygamy inequality based on the \( \alpha \)th power of concurrence, which includes the standard inequality (5) as a special case of \( \alpha = 2 \). Theorem 1 gives a weighted polygamy inequality, which can give rise to finer characterizations of entanglement distributions. Namely, once the information about the ordering of the concurrence of assistance between \( A \) and parts of the systems \( B_1B_2\cdots B_{N−1} \) are given, we can present better (weighted) characterizations of entanglement distributions among the subsystems. For example, let us consider a three-qubit state \( \rho_{ABC} \in \mathcal{H}_{ABC} \). By Theorem 1, one has \( C^\alpha(\rho_{ABC}) \leq C^\alpha(\rho_{AB}) + \alpha/2C^\alpha(\rho_{AC}) \) if \( C^\alpha(\rho_{AB}) \geq C^\alpha(\rho_{AC}) \). Since \( \alpha/2 \leq 1 \), the weight of entanglement between subsystems \( AB \) is larger than that between subsystems \( AC \). For the case that \( \alpha = 1 \), we have \( C(\rho_{ABC}) \leq C(\rho_{AB}) + 1/2C(\rho_{AC}) \). For given \( C(\rho_{ABC}) \), if \( C(\rho_{AB}) \) reduces a mount \( \Delta \), then \( C(\rho_{AC}) \) may need to increase \( 2\Delta \) to keep the inequality, which is not the case from (5), where the entanglements of subsystems \( AB \) and \( AC \) are equally weighted.

[Lemma 1.] For arbitrary two real numbers \( x, y \) such that \( x \geq y \geq 0 \), we have \( (x−y)^\alpha \geq x^\alpha − y^\alpha \) and \( (x+y)^\alpha \leq x^\alpha + y^\alpha \) for \( 0 \leq \alpha \leq 1 \).

[Proof.] \( (x−y)^\alpha \geq x^\alpha − y^\alpha \) is equivalent to \( (1−\frac{y}{x})^\alpha + (\frac{y}{x})^\alpha \geq 1 \) for nonzero \( x \). Denote \( t = \frac{y}{x} \). Then \( 0 \leq t \leq 1 \).

Set \( f(t) = (1−t)^\alpha + t^\alpha \). We have \( \frac{df}{dt} = t^{α−1} − (1−t)^{α−1} \).

When \( 0 \leq t \leq \frac{1}{2} \), then \( \frac{df}{dt} \geq 0 \), since \( t − t \geq t \) and \( α − 1 < 0 \). Therefore, \( f(t) \geq f(0) = 1 \) in this case. When \( \frac{1}{2} < t \leq 1 \), then \( \frac{df}{dt} \leq 0 \), since \( t \geq 1 − t \) and \( α − 1 \leq 0 \).

Hence, \( f(t) \geq f(1) = 1 \) in this case. In summary, for \( 0 < \alpha \leq 1 \), \( f(t)_{\min} = f(0) = f(1) = 1 \). Similarly, one can get the second inequality in Lemma 1. When \( \alpha = 0 \) or \( x = 0 \), the inequality is trivial. Hence we complete the proof of the Lemma 1.

In the following, by using the conclusion of Theorem 1 and Lemma 1, we present some monogamy inequalities and lower bounds of concurrence in terms of concurrence and COA. These monogamy relations are satisfied by the concurrence of \( N \)-qubit states under the partition \( AB \) and \( C_1\cdots C_{N−2} \), as well as under the partition \( ABC_1 \) and \( C_2\cdots C_{N−2} \), which generalize the monogamy inequalities for pure states in [24].

[Theorem 2.] For any \( 2 \otimes 2 \otimes \cdots \otimes 2 \) pure state \( ρ_{ABC\cdots C_{N−2}} \), we have

\[
C^\alpha(\rho_{AB|C_1\cdots C_{N−2}}) \\
\geq \max \left\{ \left( \sum_{i=1}^{N−2} C^2(\rho_{AC_i}) + C^2(\rho_{AB}) \right)^{\frac{\alpha}{2}} - J_B, \right. \\
\left. \left( \sum_{i=1}^{N−2} C^2(\rho_{BC_i}) + C^2(\rho_{AB}) \right)^{\frac{\alpha}{2}} - J_A \right\},
\]

for \( 0 \leq \alpha \leq 2 \), \( N \geq 4 \), where \( C^2(\rho_{AM_i}) \) is defined in (5) and \( J_A = \sum_{i=1}^{k_1} \left( \frac{c_i}{2} \right)^{2−1} C^2(\rho_{AM_i}), J_B = \sum_{i=1}^{k_2} \left( \frac{c_i}{2} \right)^{2−1} C^2(\rho_{BM_i}) \).

[Proof.] Without loss of generality, there always exists a proper ordering of the subsystems \( M_1, M_{t+1}, \cdots, M_{k_i} \) such that \( C^2(\rho_{AM_{t+i}}) \geq \sum_{i=t+1}^{k_1} C^2(\rho_{AM_i}) \) and \( C^2(\rho_{BM_{t+i}}) \geq \sum_{i=t+1}^{k_2} C^2(\rho_{BM_i}) \), \( 1 \leq t_1, t_2 \leq k − 1 \), \( 2 \leq k_1, k_2 \leq N − 1 \).

For \( 2 \otimes 2 \otimes \cdots \otimes 2 \) pure state \( ρ_{ABC\cdots C_{N−2}} \), if \( C^2(\rho_{AB|C_1\cdots C_{N−2}}) \geq C^2(\rho_{AB|AC_1\cdots C_{N−2}}) \), one has

\[
C^\alpha(\rho_{AB|C_1\cdots C_{N−2}}) \\
= (2T(\rho_{AB}))^{\frac{\alpha}{2}} \\
\geq (2T(\rho_{A}) − 2T(\rho_{B}))^{\frac{\alpha}{2}} \\
= \left[ C^2(\rho_{A|BC_1\cdots C_{N−2}}) − C^2(\rho_{B|AC_1\cdots C_{N−2}}) \right]^{\frac{\alpha}{2}} \\
\geq \left( \sum_{i=1}^{N−2} C^2(\rho_{AC_i}) + C^2(\rho_{AB}) \right)^{\frac{\alpha}{2}} - J_B,
\]

where the first inequality is due to the left inequality in [3]. Using the Lemma 1, one gets the second inequality. From (2) we get the third inequality. The last inequality is obtained by Theorem 1.

If \( C^2(\rho_{AB|C_1\cdots C_{N−2}}) \leq C^2(\rho_{B|AC_1\cdots C_{N−2}}) \), similar to the above derivation, we can obtain another inequality in Theorem 2. □
Example 2. Let us consider the following four-qubit pure state,

$$|\psi\rangle_{ABCD} = \frac{1}{\sqrt{3}}(|0000\rangle + |0010\rangle + |1011\rangle).$$  

(10)

From the Theorem 1 in [27] we get $C^2(|\psi\rangle_{AB|CD}) \geq \frac{4}{9}$ for $\alpha = 2$. From our Theorem 2, we have $C^\alpha(|\psi\rangle_{AB|CD}) \geq (\frac{3}{2})^\alpha$ for any $\alpha$, see Fig. 2.

[Remark 2]. Theorem 2 gives a monogamy inequality under the partition $AB$ and $C_1 \cdots C_{N-2}$. Similar to Theorem 1, it also gives rise finer weighted characterizations of the entanglement distributions among the subsystems, as is illustrated in Example 2. Here we indicate that the proof of the main result in Ref. [25] has a flaw: where the equality $C^2(|\psi\rangle_{A|BC}) = C^2(|\psi\rangle_{AB} + C^2(|\psi\rangle_{AC})$ is used for an arbitrary $2 \otimes 2 \otimes m$ quantum pure state $|\psi\rangle_{ABC}$, which has not been proved yet. We used some basic inequalities and Theorem 1 to prove the Theorem 2, which includes the result in Ref. [25] as a special case of $\alpha = 2$.

[Theorem 3]. For any $2 \otimes 2 \otimes \cdots \otimes 2$ pure state $|\psi\rangle_{ABC_1 \cdots C_{N-2}}$, we have

$$C^\alpha(|\psi\rangle_{AB|C_1 \cdots C_{N-2}}) \leq J_A + J_B$$  

(11)

for $0 \leq \alpha \leq 2$, $N \geq 4$, where $J_A$ and $J_B$ are defined the same with Theorem 2.

[Proof]. Without loss of generality, there always exists a proper ordering of the subsystems $M_1, M_{t_1+1}, \cdots, M_{t_2}$ ($i = 1, 2$) such that $C^2_a(\rho_{M_1}) \geq \sum_{t_1+1}^{k_1} C^2_a(\rho_{M_2})$ and $C^2_a(\rho_{B|M_2}) \geq \sum_{t_2+1}^{k_2} C^2_a(\rho_{B|M_1})$, $1 \leq t_1, t_2 \leq k - 1, 2 \leq k_1, k_2 \leq N - 1$.

For $2 \otimes 2 \otimes \cdots \otimes 2$ pure state $|\psi\rangle_{ABC_1 \cdots C_{N-2}}$, one has

$$C^\alpha(|\psi\rangle_{AB|C_1 \cdots C_{N-2}})$$

$$= (2T(\rho_{AB}))^\frac{\alpha}{2}$$

$$\leq (2T(\rho_A) + 2T(\rho_B))^\frac{\alpha}{2}$$

$$= (C^2(|\psi\rangle_{A|BC_1 \cdots C_{N-2}}) + C^2(\rho_{B|A|C_1 \cdots C_{N-2}}))^\frac{\alpha}{2}$$

$$\leq C^\alpha(|\psi\rangle_{A|BC_1 \cdots C_{N-2}}) + C^\alpha(\rho_{B|A|C_1 \cdots C_{N-2}})$$

$$\leq J_A + J_B,$$

where the first inequality is due to the right inequality in (3). The second inequality is due to Lemma 1. Using the Theorem 1, one gets the last inequality. □

For the four-qubit state $|00\rangle$, we have $C_\alpha(\rho_{AB}) = 0$. Then from the Theorem 2 in [27], we get $C^2(|\psi\rangle_{AB|CD}) \leq \frac{4}{3}$ for $\alpha = 2$. From our Theorem 3, we have $C^\alpha(|\psi\rangle_{AB|CD}) \leq (\frac{2}{3})^\alpha + \frac{2}{3}(\frac{3}{2})^\alpha$ for any $\alpha$, see Fig. 3.

[Remark 3]. We have presented a weighted polygamy inequality for the $\alpha$th ($0 \leq \alpha \leq 2$) power of concurrence, which includes the result in Ref. [25] as a special case of $\alpha = 2$. For $0 \leq \alpha < 2$, Theorem 3 gives rise to again finer characterizations of the entanglement distributions by weighting the entanglement of the subsystems.

Now we generalize our results to the concurrence $C_{ABC_1|C_2 \cdots C_{N-2}}(\rho_{AB})$ under the partition $ABC_1$ and $C_2 \cdots C_{N-2}$ ($N \geq 6$) for pure state $|\psi\rangle_{ABC_1 \cdots C_{N-2}}$. Similar to Theorem 2 and Theorem 3, we obtain the following corollaries:

[Corollary 1]. For any $N$-qubit pure state $|\psi\rangle_{ABC_1 \cdots C_{N-2}}$, we have

$$C^\alpha(|\psi\rangle_{ABC_1|C_2 \cdots C_{N-2}})$$

$$\geq \max \left\{ \left( \sum_{i=1}^{N-2} C^2(\rho_{AC_i}) + C^2(\rho_{AB}) \right)^\frac{\alpha}{2} - J_B, \left( \sum_{i=1}^{N-2} C^2(\rho_{B|C_i}) + C^2(\rho_{AB}) \right)^\frac{\alpha}{2} - J_A \right\} - J_{C_1},$$

(12)

where $J_A$, $J_B$ are defined as in Theorem 2, $J_{C_1} = \sum_{i=1}^{k_1} (\frac{3}{2})^{i-1} C^2_a(\rho_{C_1|M_i})$, $2 \leq m \leq N - 3$, $N \geq 6$.

[Proof]. For any $N$-qubit pure state $|\psi\rangle_{ABC_1 \cdots C_{N-2}}$, if
Namely, bound (12) is better than (13) in this case.

Combining Theorem 1 and Theorem 2, we obtain (12).

[Corollary 2]. For any N-qubit pure state $|\psi\rangle_{ABC_1\cdots C_{N-2}}$, if $C^2(|\psi\rangle_{AB|C_1\cdots C_{N-2}}) \leq C^2(|\psi\rangle_{C_1ABC_2\cdots C_{N-2}})$, we have

\[
C^a(|\psi\rangle_{ABC_1|C_2\cdots C_{N-2}}) \\
\geq (2T(\rho_{ABC_1}))^{\frac{2}{2}} \geq |2T(\rho_{AB}) - 2T(\rho_{C_1})|^{\frac{2}{2}} = |C^2(|\psi\rangle_{AB|C_1\cdots C_{N-2}}) - C^2(|\psi\rangle_{C_1ABC_2\cdots C_{N-2}})|^{\frac{2}{2}} \geq C^a(|\psi\rangle_{AB|C_1\cdots C_{N-2}}) - C^o(|\psi\rangle_{C_1ABC_2\cdots C_{N-2}}),
\]

where the first inequality is due to $T(\rho_{ABC_1}) \geq T(\rho_{AB}) - T(\rho_{C_1})$. Using Lemma 1, we get the second inequality. Combining Theorem 1 and Theorem 2, we obtain (12).

Example 4. For the state $|\psi\rangle_{ABC_1C_2C_3C_4} = \frac{1}{\sqrt{2}}(|000000\rangle + |010000\rangle)$, we have $C(\rho_{AB}) = C(\rho_{ABC_1}) = C(\rho_{AC_1C_3}) = C(\rho_{C_1C_2C_3}) = 0$ for $i = 2, 3, 4$; $C(\rho_{BC_1}) = C(\rho_{BC_2}) = 0$ for $i = 1, 2, 3, 4$; and $C(\rho_{C_1C_3}) = C(\rho_{C_1C_4}) = 1$. Hence, we have $C(|\psi\rangle) \geq 1$ from (12) and $C(|\psi\rangle) \geq 0$ from (13). Namely, bound (12) is better than (13) in this case.

Example 5. For the state $|\psi\rangle_{ABC_1C_2C_3C_4} = \frac{1}{\sqrt{2}}(|000000\rangle + |001100\rangle)$, it is straightforward to calculate that $C(\rho_{AB}) = C(\rho_{AC_1}) = C(\rho_{BC_1}) = C(\rho_{BC_2}) = C(\rho_{C_1C_3}) = C(\rho_{C_1C_4}) = 0$ for $i = 2, 3, 4$; and $C(\rho_{C_1C_2}) = C(\rho_{C_2C_3}) = 1$. Hence, we have $C(|\psi\rangle) \geq 0$ from (12) and $C(|\psi\rangle) \geq 1$ from (13). The bound (13) is better than (12) in this case.

GENERALIZED MONOGAMY AND POLYGAMY

RELATIONS OF NEGATIVITY

Another well-known quantifier of bipartite entanglement is the negativity. Given a bipartite state $\rho_{AB}$ in $H_A \otimes H_B$, the negativity is defined by

\[
N(\rho_{AB}) = \frac{1}{2} ||\rho_{AB}^T||_1 - 1,
\]

where $\rho_{AB}^T$ is the partially transposed matrix of $\rho_{AB}$ with respect to the subsystem $A$. $||X||$ denotes the trace norm of $X$, i.e. $||X|| = \text{Tr}\sqrt{XX^\dagger}$. Negativity is a computable measure of entanglement, and is a convex function of $\rho_{AB}$. It vanishes if and only if $\rho_{AB}$ is separable for the 2$\otimes$ 2 and 2$\otimes$ 3 systems. For the purposes of discussion, we use the following definition of negativity:

\[
N(\rho_{AB}) = ||\rho_{AB}^T||_1 - 1.
\]

For any bipartite pure state $|\psi\rangle_{AB}$ in a $d \otimes d$ quantum system with Schmidt rank $d$, $|\psi\rangle_{AB} = \sum_{i=1}^{d} \sqrt{\lambda_i} |ii\rangle$, one has

\[
N(|\psi\rangle_{AB}) = 2 \sum_{i<j} \sqrt{\lambda_i \lambda_j},
\]

from the definition of concurrence, we have

\[
C(|\psi\rangle_{AB}) = 2 \sqrt{\sum_{i<j} \lambda_i \lambda_j}.
\]

Combining (15) with (16), one obtains

\[
N(|\psi\rangle_{AB}) \geq C(|\psi\rangle_{AB}).
\]

For any bipartite pure state $|\psi\rangle_{AB}$ with Schmidt rank 2, one has $N(|\psi\rangle_{AB}) = C(|\psi\rangle_{AB})$ from (15) and (16). For a mixed state $\rho_{AB}$, the convex-roof extended negativity (CREN) is defined by

\[
N_c(\rho_{AB}) = \min_i p_i N(|\psi_i\rangle_{AB}),
\]

where the minimum is taken over all possible pure state decompositions $\{p_i, |\psi_i\rangle_{AB}\}$ of $\rho_{AB}$. CREN gives a perfect discrimination of positively partial transposed bound entangled states and separable states in any bipartite quantum systems. For a mixed state $\rho_{AB}$, the convex-roof extended negativity of assistance (CRENOA) is defined by

\[
N_a(\rho_{AB}) = \max_i p_i N(|\psi_i\rangle_{AB}),
\]

where the maximum is taken over all possible pure state decompositions $\{p_i, |\psi_i\rangle_{AB}\}$ of $\rho_{AB}$.
CREN is equivalent to concurrence for any pure state with Schmidt rank 2 \[ N \leq \] for 0 \[ \leq \] M \[ | \) from [17], [18], [19] and the monogamy of the concurrence, we have
\[ N^\alpha(\psi_{AB|B_1B_2\ldots B_{N-1}}) \geq N^\alpha_c(\rho_{AB_1}) + N^\alpha_c(\rho_{AB_2}) + \cdots + N^\alpha_c(\rho_{AB_{N-1}}) \]
for \( \alpha \geq 2 \). The dual inequality \[ N^2(\psi_{AB|B_1B_2\ldots B_{N-1}}) \]
\[ \leq N^2(\rho_{AB_1}) + N^2(\rho_{AB_2}) + \cdots + N^2(\rho_{AB_{N-1}}) \]
(21)
where \( N^2(\rho_{AM_i}) = \sum_{j=M_{i-1}+1}^{M_i} N^2(\rho_{AB}) \) with \( M_0 = 0 \), \( \sum_{i=1}^{k} M_i = N - 1 \), 1 \( \leq k \leq N - 1 \). By similar consideration to concurrence, we get the upper bound of the \( \alpha \)-th power of negativity as follows.

**Theorem 4.** For any \( N \)-qubit pure state \( |\psi_{AB|B_1B_2\ldots B_{N-1}}\rangle \), we have
\[ N^\alpha(\psi_{AB|B_1B_2\ldots B_{N-1}}) \]
\[ \leq N^\alpha(\rho_{AM_1}) + \alpha^2 N^\alpha(\rho_{AM_2}) + \cdots + \binom{k}{2}^{-1} N^\alpha(\rho_{AM_k}) \]
(22)
for all 0 \( \leq \alpha \leq 2 \), \( N \geq 4 \).

**Theorem 5.** For any \( 2 \otimes 2 \otimes \cdots \otimes 2 \) pure state \( |\psi_{ABC_1\ldots C_{N-2}}\rangle \), we have
\[ N^\alpha(|\psi_{ABC|C_1\ldots C_{N-2}}\rangle) \]
\[ \geq \max\left(\sum_{i=1}^{N-2} N^2(\rho_{AC_i}) + N^2(\rho_{AB})\right)^{\frac{\alpha}{2}} - J_B', \]
\[ \left(\sum_{i=1}^{N-2} N^2(\rho_{BC_i}) + N^2(\rho_{AB})\right)^{\frac{\alpha}{2}} - J_A' \]
for 0 \( < \alpha \leq 2 \), \( N \geq 4 \), where \( J_A' = \sum_{i=1}^{k_1} \left(\frac{\alpha}{2}\right)^{i-1} N^\alpha(\rho_{AM_i}) \), \( J_B' = \sum_{i=1}^{k_2} \left(\frac{\alpha}{2}\right)^{i-1} N^\alpha(\rho_{BM_i}) \).

**Proof.** Without loss of generality, there always exists a proper ordering of the subsystems \( M_1, M_{t_1+1}, \ldots, M_{t_1} \) (i = 1, 2) such that \( N^2(\rho_{AM_1}) \geq \sum_{i=t_1+1}^{k_1} N^2(\rho_{AM_i}) \) and
\[ N^2(\rho_{BM_1}) \geq \sum_{i=t_2+1}^{k_2} N^2(\rho_{BM_i}), 1 \leq t_1, t_2 \leq k - 1, 2 \leq k_1, k_2 \leq N - 1. \]

For \( 2 \otimes 2 \otimes \cdots \otimes 2 \) pure state \( |\psi_{ABC_1\ldots C_{N-2}}\rangle \), we have
\[ N^\alpha(\psi_{AB|C_1\ldots C_{N-2}}) \geq \max\left(\sum_{i=1}^{N-2} C^2(\rho_{AC_i}) + C^2(\rho_{AB})\right)^{\frac{\alpha}{2}} - J_B, \]
\[ \left(\sum_{i=1}^{N-2} C^2(\rho_{BC_i}) + C^2(\rho_{AB})\right)^{\frac{\alpha}{2}} - J_A \]
\[ = \max\left(\sum_{i=1}^{N-2} N^2(\rho_{AC_i}) + N^2(\rho_{AB})\right)^{\frac{\alpha}{2}} - J_B', \]
\[ \left(\sum_{i=1}^{N-2} N^2(\rho_{BC_i}) + N^2(\rho_{AB})\right)^{\frac{\alpha}{2}} - J_A', \]
where the first inequality is due to (17), the second inequality is from Theorem 2, the equality is based on (18) and (19).

For \( N \)-qubit pure state \( |\psi_{AB|B_1B_2\ldots B_{N-1}}\rangle \), based on the result in [33] [34], \( N(|\psi_{AB|C_1\ldots C_{N-2}}\rangle) \leq \sqrt{\frac{r^2-1}{2}C(\psi_{AB|C_1\ldots C_{N-2}})} \), where \( r \) is the Schmidt rank of the pure state \( |\psi_{ABC_1\ldots C_{N-2}}\rangle \). From Theorem 3, we can obtain the upper bound of negativity under the partition \( AB \) and \( C_1 \cdots C_{N-2} \).

**Theorem 6.** For any \( 2 \otimes 2 \otimes \cdots \otimes 2 \) pure state \( |\psi_{ABC_1\ldots C_{N-2}}\rangle \), we have
\[ N^\alpha(\psi_{AB|C_1\ldots C_{N-2}}) \leq \left(\frac{r(r-1)}{2}\right)^{\frac{\alpha}{2}} (J_A' + J_B'), \]
for 0 \( \leq \alpha \leq 2 \), where \( J_A', J_B' \) are given in Theorem 5.

**Example 5.** Let us consider the 4-qubit generalized W-class state,
\[ |W\rangle_{ABC_1C_2} = \lambda_1|0000\rangle + \lambda_2|0100\rangle \]
\[ + \lambda_3|0011\rangle + \lambda_4|0101\rangle \]
(23)
where \( \sum_1^4 \lambda_i^2 = 1 \). We have \( N(|W\rangle_{ABC_1C_2}) = 2\sqrt{(\lambda_1^2 + \lambda_2^2)(\lambda_3^2 + \lambda_4^2)} \), \( N_c(\rho_{AB}) = N_c(\rho_{AB_2}) = 2\lambda_1\lambda_2 \), \( N_c(\rho_{AC_1}) = N_c(\rho_{AC_2}) = 2\lambda_1\lambda_3 \), \( N_c(\rho_{AC_3}) = 2\lambda_1\lambda_4 \). Taking \( \lambda_1 = \frac{1}{4}, \lambda_2 = \frac{1}{2}, \lambda_3 = \frac{\sqrt{2}}{2}, \lambda_4 = \frac{1}{4} \), we get \( J_A' = J_B' = \left(\frac{4}{5}\right)^{\alpha} + \frac{\alpha}{2} \left(\frac{4}{5}\right)^{\alpha} \left(\frac{4}{5}\right)^{\alpha} \). Set \( y = N^\alpha(|W\rangle_{ABC_1C_2}) - ((N^2_c(\rho_{AB}) + N^2_c(\rho_{AC_1}) + N^2_c(\rho_{AC_2}))^{\frac{\alpha}{2}} - J_A') \), i.e., the difference between the left and right side of Theorem 5. We elucidate the results of Theorem 5 in Fig. 4 for this case. For the case that the Schmidt rank of the pure state of \( |\psi_{ABC_1\ldots C_{N-2}}\rangle \) is 2, we show the results of Theorem 6 in Fig. 5.

**Conclusion**

Entanglement monogamy and polygamy are fundamental properties of multipartite entangled
states. We have presented the monogamy relations of the $\alpha$th power of concurrence for $N$-qubit systems by showing the relations among $C(|\psi\rangle_{AB}|C_1\cdots C_{N-2})$, $C(\rho_{AB})$, $C(\rho_{AC})$, $C(\rho_{BC})$, $C_0(\rho_{AC})$, and $C_0(\rho_{BC})$, $1 \leq i \leq N-2$, which give rise to the lower and upper bounds on the entanglement sharing among the partitions. We have investigated the monogamy relations based on concurrence and COA. The upper bound of the $\alpha$th ($0 \leq \alpha \leq 2$) power of concurrence has been obtained based on the COA. We have obtained the monogamy and polygamy relations satisfied by the $\alpha$th ($0 \leq \alpha \leq 2$) power of concurrence in $N$-qubit pure states under the partition $AB$ and $C_1\cdots C_{N-2}$, as well as under the partition $ABC_1$ and $C_1\cdots C_{N-2}$. These relations also give rise to a kind of trade-off relations related to the lower and upper bounds of concurrences. Based on the relations between negativity and concurrence, we have also obtained the similar results for CREN and CRENOA. These results may be generalized to monogamy and polygamy relations under arbitrary partitions $C_{ABC_1\cdots C_i|C_{i+1}\cdots C_{N-2}}$, $2 \leq i \leq N-2$. Our approach may be used for the investigation of entanglement distribution based on other measures of quantum correlations.

Acknowledgments This work is supported by the NSF of China under Grant No. 11675113 and NSF of Beijing under No. KZ201810028042.

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