COCO: The Large Scale Black-Box Optimization Benchmarking (bbob-largescale) Test Suite

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Abstract

The bbob-largescale test suite, containing 24 single-objective functions in continuous domain, extends the well-known single-objective noiseless bbob test suite [HAN2009], which has been used since 2009 in the BBOB workshop series, to large dimension. The core idea is to make the rotational transformations $R, Q$ in search space that appear in the bbob test suite computationally cheaper while retaining some desired properties. This documentation presents an approach that replaces a full rotational transformation with a combination of a block-diagonal matrix and two permutation matrices in order to construct test functions whose computational and memory costs scale linearly in the dimension of the problem.

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1 Introduction

In the \texttt{bbob-largescale} test suite, we consider single-objective, unconstrained minimization problems of the form

$$\min_{x \in \mathbb{R}^n} f(x),$$

with problem dimensions $n \in \{20, 40, 80, 160, 320, 640\}$.

The objective is to find, as quickly as possible, one or several solutions $x$ in the search space $\mathbb{R}^n$ with \textit{small} value(s) of $f(x) \in \mathbb{R}$. We generally measure the \textit{time} of an optimization run as the number of calls to (queries of) the objective function $f$.

We remind in the next sections some notations and definitions.

1.1 Terminology

\textbf{Function} We talk about an objective \textit{function} $f$ as a parametrized mapping $\mathbb{R}^n \rightarrow \mathbb{R}$ with scalable input space, that is, $n$ is not (yet) determined. Functions are parametrized such that different \textit{instances} of the “same” function are available, e.g. translated or rotated versions.

\textbf{Problem} We talk about a \textit{problem}, \texttt{coco_problem_t}, as a specific \textit{function instance} on which an optimization algorithm is run. Specifically, a problem can be described as the triple \texttt{(dimension, function, instance)}. A problem can be evaluated and returns an $f$-value. In the context of performance assessment, a target $f$- or indicator-value is attached to each problem. That is, a target value is added to the above triple to define a single problem in this case.

\textbf{Runtime} We define \textit{runtime}, or \textit{run-length} as the \textit{number of evaluations} conducted on a given problem, also referred to as number of \textit{function} evaluations. Our central performance measure is the runtime until a given target value is hit.

\textbf{Suite} A test- or benchmark-suite is a collection of problems, typically between twenty and a hundred.

1.2 Functions, Instances and Problems

Each function is \textit{parametrized} by the (input) dimension, $n$, its identifier $i$, and the instance number, $j$, that is:

$$f^j_i \equiv f(n, i, j) : \mathbb{R}^n \rightarrow \mathbb{R}, \quad x \mapsto f^j_i(x) = f(n, i, j)(x).$$

Varying $n$ or $j$ leads to a variation of the same function $i$ of a given suite. By fixing $n$ and $j$ for function $f_i$, we define an optimization \textbf{problem} $(n, i, j) \equiv (f_i, n, j)$ that can be presented to the optimization algorithm. Each problem receives again an index in the suite, mapping the triple $(n, i, j)$ to a single number.
We can think of \( j \) as an index to a continuous parameter vector setting, as it parametrizes, among others things, translations and rotations. In practice, \( j \) is the discrete identifier for single instantiations of these parameters.

### 1.3 Runtime and Target Values

In order to measure the runtime of an algorithm on a problem, we establish a hitting time condition. For a given problem \((f, n, j)\), we prescribe a **target value** \( t \) as a specific \( f \)-value of interest \[\text{HAN2016perf}\]. For a single run, when an algorithm reaches or surpasses the target value \( t \) on problem \((f, n, j)\), we say that it has solved the problem \((f, n, j, t)\) — it was successful.\(^1\)

The **runtime** is, then, the evaluation count when the target value \( t \) was reached or surpassed for the first time. That is, the runtime is the number of \( f \)-evaluations needed to solve the problem \((f, n, j, t)\).\(^2\) **Measured runtimes are the only way how we assess the performance of an algorithm.** Observed success rates are generally translated into runtimes on a subset of problems.

If an algorithm does not hit the target in a single run, its runtime remains undefined — while, then, this runtime is bounded from below by the number of evaluations in this unsuccessful run. The number of available runtime values depends on the budget the algorithm has explored (the larger the budget, the more likely the target-values are reached). Therefore, larger budgets are preferable — however they should not come at the expense of abandoning reasonable termination conditions. Instead, restarts should be done \[\text{HAN2016ex}\].

### 2 Overview of the Proposed bbob-largescale Test Suite

The **bbob-largescale** test suite provides 24 functions in six dimensions (20, 40, 80, 160, 320 and 640) within the COCO framework \[\text{HAN2016co}\]. It is derived from the existing single-objective, unconstrained bbob test suite with modifications that allow the user to benchmark algorithms on high dimensional problems efficiently. We will explain in this section how the bbob-largescale test suite is built.

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\(^1\) Note the use of the term **problem** in two meanings: as the problem the algorithm is benchmarked on, \((f, n, j)\), and as the problem, \((f, n, j, t)\), an algorithm can solve by hitting the target \( t \) with the runtime, \(\text{RT}(f, n, j, t)\), or may fail to solve. Each problem \((f, n, j)\) gives raise to a collection of dependent problems \((f, n, j, t)\). Viewed as random variables, the events \(\text{RT}(f, n, j, t)\) given \((f, n, j)\) are not independent events for different values of \( t \).

\(^2\) Target values are directly linked to a problem, leaving the burden to properly define the targets with the designer of the benchmark suite. The alternative is to present final \( f \)-values as results, leaving the (rather unsurmountable) burden to interpret these values to the reader. Fortunately, there is an automatized generic way to generate target values from observed runtimes, the so-called run-length based target values \[\text{HAN2016perf}\].
2.1 The single-objective bbob functions

The bbob test suite relies on the use of a number of raw functions from which 24 bbob functions are generated. Initially, so-called raw functions are designed. Then, a series of transformations on these raw functions, such as linear transformations (e.g., translation, rotation, scaling) and/or non-linear transformations (e.g., $T_{osz}, T_{asy}$) will be applied to obtain the actual bbob test functions. For example, the test function $f_{13}(x)$ (Sharp Ridge function) with (vector) variable $x$ is derived from a raw function defined as follows:

$$f_{\text{Sharp Ridge}}(z) = z_1^2 + 100 \sum_{i=2}^{n} z_i^2.$$  

Then one applies a sequence of transformations: a translation by using the vector $x^\text{opt}$; then a rotational transformation $R$; then a scaling transformation $\Lambda^{10}$; then another rotational transformation $Q$ to get the relationship $z = Q \Lambda^{10} R (x - x^\text{opt})$; and finally a translation in objective space by using $f^\text{opt}$ to obtain the final function in the testbed:

$$f_{13}(x) = f_{\text{Sharp Ridge}}(z) + f^\text{opt}.$$  

There are two main reasons behind the use of transformations here:

1. provide non-trivial problems that cannot be solved by simply exploiting some of their properties (separability, optimum at fixed position, ...)
2. allow to generate different instances, ideally of similar difficulty, of the same problem by using different (pseudo-)random transformations.

Rotational transformations are used to avoid separability and thus coordinate system dependence in the test functions. The rotational transformations consist in applying an orthogonal matrix to the search space: $x \rightarrow z = Rx$, where $R$ is the orthogonal matrix. While the other transformations used in the bbob test suite could be naturally extended to the large scale setting due to their linear complexity, rotational transformations have quadratic time and space complexities. Thus, we need to reduce the complexity of these transformations in order for them to be usable, in practice, in the large scale setting.

2.2 Extension to large scale setting

Our objective is to construct a large scale test suite where the cost of a function call is acceptable in higher dimensions while preserving the main characteristics of the original functions in the bbob test suite. To this end, we will replace the full orthogonal matrices of the rotational transformations, which would be too expensive in our large scale setting, with orthogonal transformations that have linear complexity in the problem dimension: permuted orthogonal block-diagonal matrices ([AIT2016]).

Specifically, the matrix of a rotational transformation $R$ will be represented as:

$$R = P_{\text{left}}BP_{\text{right}}.$$
Here, $P_{\text{left}}$ and $P_{\text{right}}$ are two permutation matrices\(^3\) and $B$ is a block-diagonal matrix of the form:

$$B = \begin{pmatrix} B_1 & 0 & \ldots & 0 \\ 0 & B_2 & \ldots & 0 \\ 0 & 0 & \ddots & 0 \\ 0 & 0 & \ldots & B_{n_b} \end{pmatrix},$$

where $n_b$ is the number of blocks and $B_i, 1 \leq i \leq n_b$ are square matrices of sizes $s_i \times s_i$ satisfying $s_i \geq 1$ and $\sum_{i=1}^{n_b} s_i = n$. In this case, the matrices $B_i, 1 \leq i \leq n_b$ are all orthogonal. Thus, the matrix $B$ is also an orthogonal matrix.

This representation allows the rotational transformation $R$ to satisfy three desired properties:

1. Have (almost) linear cost (due to the block structure of $B$).
2. Introduce non-separability.
3. Preserve the eigenvalues and therefore the condition number of the original function when it is convex quadratic (since $R$ is orthogonal).

### 2.3 Generating the orthogonal block matrix $B$

The block-matrices $B_i, i = 1, 2, \ldots, n_b$ will be uniformly distributed in the set of orthogonal matrices of the same size. To this end, we first generate square matrices with sizes $s_i (i=1,2,\ldots,n_b)$ whose entries are i.i.d. standard normally distributed. Then we apply the Gram-Schmidt process to orthogonalize these matrices.

The parameters of this procedure include:

- the dimension of the problem $n$,
- the block sizes $s_1, \ldots, s_{n_b}$, where $n_b$ is the number of blocks. In this test suite, we set $s_i = s := \min\{n, 40\} \forall i = 1, 2, \ldots, n_b$ (except, maybe, for the last block which can be smaller)\(^4\) and thus $n_b = \lceil n/s \rceil$.

### 2.4 Generating the permutation matrices $P$

In order to generate the permutation matrix $P$, we start from the identity matrix and apply, successively, a set of so-called truncated uniform swaps. Each row/column (up to a maximum number of swaps) is swapped with a row/column chosen uniformly from the set of rows/columns within a fixed range $r_s$. A random order of the rows/columns is generated to avoid biases towards the first rows/columns.

\(^3\)A permutation matrix is a square binary matrix that has exactly one entry of 1 in each row and each column and 0s elsewhere.

\(^4\)This setting allows to have the problems in dimensions 20 and 40 overlap between the bbob test suite and its large-scale extension since in these dimensions, the block sizes coincide with the problem dimensions.
Let $i$ be the index of the first variable/row/column to be swapped and $j$ be the index of the second swap variable. Then

$$j \sim U(\{l_b(i), l_b(i) + 1, \ldots, u_b(i)\} \backslash \{i\}),$$

where $U(S)$ is the uniform distribution over the set $S$ and $l_b(i) = \max(1, i - r_s)$ and $l_b(i) = \min(n, i + r_s)$ with $r_s$ a parameter of the approach. If $r_s \leq (n - 1)/2$, the average distance between the first and the second swap variable ranges from $(\sqrt{2} - 1)r_s + 1/2$ (in the case of an asymmetric choice for $j$, i.e. when $i$ is chosen closer to 1 or $n$ than $r_s$) to $r_s/2 + 1/2$ (in the case of a symmetric choice for $j$). It is maximal when the first swap variable is at least $r_s$ away from both extremes or is one of them.

**Algorithm 1** below describes the process of generating a permutation using a series of truncated uniform swaps with the following parameters:

- $n$, the number of variables,
- $n_s$, the number of swaps,
- $r_s$, the swap range.

Starting with the identity permutation $p$ and another permutation $\pi$, drawn uniform at random, we apply the swaps defined above by taking $p_{\pi}(1), p_{\pi}(2), \ldots, p_{\pi}(n_s)$, successively, as first swap variable. The resulting vector $p$ will be the desired permutation.

**Algorithm 1: Truncated Uniform Permutations**

```
1. $p \leftarrow (1, \ldots, n)$
2. Generate a permutation $\pi$ uniformly at random
3. for $1 \leq k \leq n_s$ do
   4. • $i \leftarrow \pi(k)$, i.e., $p_{\pi(k)}$ is the first swap variable
   5. • $l_b \leftarrow \max(1, i - r_s)$
   6. • $u_b \leftarrow \min(n, i + r_s)$
   7. • $S \leftarrow \{l_b, l_b + 1, \ldots, u_b\} \backslash \{i\}$
   8. • Sample $j$ uniformly at random in $S$
   9. • Swap $p_i$ and $p_j$
4. end for
11. return $p$
```

In this test suite, we set $n_s = n$ and $r_s = \lfloor n/3 \rfloor$. Some numerical results in [AIT2016] show that with such parameters, the proportion of variables that are moved from their original position when
applying Algorithm 1 is approximately 100% for all dimensions 20, 40, 80, 160, 320, and 640 of
the bbob-largescale test suite.

2.5 Implementation

Now, we describe how these changes to the rotational transformations are implemented with the
realizations of $P_{\text{left}}B_{\text{right}}$. This will be illustrated through an example on the Ellipsoidal function
(rotated) $f_{10}(x)$ (see the table in the next section), which is defined by

$$f_{10}(x) = \gamma(n) \times \sum_{i=1}^{n} 10^{6i-1} z_{i}^{2} + f_{\text{opt}}, \text{with } z = T_{\text{osz}}(R(x - x_{\text{opt}})), R = P_{1}BP_{2},$$

as follows:

(i) First, we obtain the three matrices needed for the transformation, $B, P_{1}, P_{2}$, as follows:

```python
    coco_compute_blockrotation(B, seed1, n, s, n_b);
coco_compute_truncated_uniform_swap_permutation(P1, seed2, n, n_s, r_s);
coco_compute_truncated_uniform_swap_permutation(P2, seed3, n, n_s, r_s);
```

2. Then, wherever in the bbob test suite, we use the following

```python
    problem = transform_vars_affine(problem, R, b, n);
```

to make a rotational transformation, then in the bbob-largescale test suite, we replace it with the three transformations

```python
    problem = transform_vars_permutation(problem, P2, n);
    problem = transform_vars_blockrotation(problem, B, n, s, n_b);
    problem = transform_vars_permutation(problem, P1, n);
```

Here, $n$ is again the problem dimension, $s$ the size of the blocks in $B$, $n_{b}$ : the number of blocks,
$n_{s}$ : the number of swaps, and $r_{s}$ : the swap range as presented previously.

**Important remark:** Although the complexity of bbob test suite is reduced considerably by the
above replacement of rotational transformations, we recommend running the experiment on the
bbob-largescale test suite in parallel.

3 Functions in bbob-largescale test suite

The table below presents the definition of all 24 functions of the bbob-largescale test suite in
detail. Beside the important modification on rotational transformations, we also make two changes
to the raw functions in the bbob test suite.
• All functions, except for the Schwefel, Schaffer, Weierstrass, Gallagher, and Katsuura functions, are normalized by the parameter $\gamma(n) = \min(1, 40/n)$ to have uniform target values that are comparable, in difficulty, over a wide range of dimensions.

• The Discus, Bent Cigar and Sharp Ridge functions are generalized such that they have a constant proportion of distinct axes that remain consistent with the bbob test suite.

For a better understanding of the properties of these functions and for the definitions of the used transformations and abbreviations, we refer the reader to the original bbob function documentation for details.

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Table 1: Function descriptions of the separable, moderate, and ill-conditioned function groups of the \texttt{bbob-largescale} test suite.

| Group 1: Separable functions | Formulation | Transformations |
|------------------------------|-------------|-----------------|
| Sphere Function              | $f_1(x) = \gamma(n) \times \sum_{i=1}^n z_i^2 + f_{\text{opt}}$ | $z = x - x^{\text{opt}}$ |
| Ellipsoidal Function         | $f_2(x) = \gamma(n) \times \sum_{i=1}^{n} 10^{\frac{4}{1-n}} z_i^2 + f_{\text{opt}}$ | $z = T_{\text{ell}}(x - x^{\text{opt}})$ |
| Rastrigin Function           | $f_3(x) = \gamma(n) \times (10n - 10) \sum_{i=1}^n \cos(2\pi z_i) + ||z||^2 + f_{\text{opt}}$ | $z = A_{10}^{\text{opt}}(T_{\text{ell}}(x - x^{\text{opt}}))$ |
| Bueche-Rastrigin Function    | $f_4(x) = \gamma(n) \times (10n - 10) \sum_{i=1}^n \cos(2\pi z_i) + ||z||^2 + 100f_{\text{pen}}(x) + f_{\text{opt}}$ | $z_i = s_i T_{\text{ell}}(x_i - x_i^{\text{opt}})$ for $i = 1, \ldots, n$
| Linear Slope                 | $f_5(x) = \gamma(n) \times \sum_{i=1}^n (5s_i) - s_i z_i + f_{\text{opt}}$ | $z_i = \begin{cases} 
10 \times 10^{\frac{2}{1-\frac{n}{4}}} & \text{if } x_i > 0 \text{ and } i \text{ odd} \\
10^{\frac{2}{1-\frac{n}{4}}} & \text{otherwise}
\end{cases}$ for $i = 1, \ldots, n$ |

| Group 2: Functions with low or moderate conditioning | | |
| Attractive Sector Function   | $f_6(x) = T_{\text{ell}} \left( \gamma(n) \times \sum_{i=1}^n (s_i z_i)^2 \right)^{1/2} + f_{\text{opt}}$ | $z = Q A_{10}^{\text{opt}} R(x - x^{\text{opt}})$ with $R = P_{11} B_{11} P_{12}$, $Q = P_{21} B_{21} P_{22}$, $s_i = 10^{2^j}$ if $x_i \times x_i^{\text{opt}} > 0$ for $i = 1, \ldots, n$
| Step Ellipsoidal Function    | $f_7(x) = \gamma(n) \times 0.1 \max \left( \frac{1}{10^4}, \frac{10^2}{1-n} \sum_{i=1}^n z_i^2 \right) + f_{\text{pen}}(x) + f_{\text{opt}}$ | $z = \begin{cases} 
\text{otherwise}
\end{cases}$ for $i = 1, \ldots, n$ |
| Rosenbloc Function, original | $f_8(x) = \gamma(n) \times \sum_{i=1}^{\lfloor n/4 \rfloor} \left( 100 (z_i^2 - z_{i+1}^2)^2 + (z_i - 1)^2 \right) + f_{\text{opt}}$ | $z = max \left( 1, \frac{\sqrt{8}}{8} \right) (x - x^{\text{opt}}) + 1$, $x^{\text{opt}} \in [-3,3]^n$ |
| Rosenbloc Function, rotated  | $f_9(x) = \gamma(n) \times \sum_{i=1}^{\lfloor n/4 \rfloor} \left( 100 (z_i^2 - z_{i+1}^2)^2 + (z_i - 1)^2 \right) + f_{\text{opt}}$ | $z = max \left( 1, \frac{\sqrt{8}}{8} \right) R(x - x^{\text{opt}}) + 1$ with $R = P_{11} B_{12} P_{22}$, $x^{\text{opt}} \in [-3,3]^n$ |

| Group 3: Functions with high conditioning and unimodal | | |
| Ellipsoidal Function         | $f_{10}(x) = \gamma(n) \times \sum_{i=1}^n 10^{\frac{4}{1-n}} z_i^2 + f_{\text{opt}}$ | $z = T_{\text{ell}}(R(x - x^{\text{opt}}))$ with $R = P_{11} B_{21}$ |
| Discus Function              | $f_{11}(x) = \gamma(n) \times \left( 10^6 \sum_{i=1}^{\lfloor n/40 \rfloor} z_i^2 + \sum_{i=\lfloor n/40 \rfloor + 1}^n z_i^2 \right) + f_{\text{opt}}$ | $z = T_{\text{ell}}(R(x - x^{\text{opt}}))$ with $R = P_{11} B_{22}$ |
| Bent Cigar Function          | $f_{12}(x) = \gamma(n) \times \sum_{i=1}^{\lfloor n/40 \rfloor} z_i^2 + 10^6 \sum_{i=\lfloor n/40 \rfloor + 1}^n z_i^2 + f_{\text{opt}}$ | $z = R T_{\text{opt}}^{10} (R(x - x^{\text{opt}}))$ with $R = P_{11} B_{11}$ |
| Sharp Ridge Function         | $f_{13}(x) = \gamma(n) \times \sum_{i=1}^{\lfloor n/40 \rfloor} z_i^2 + 100 \sum_{i=\lfloor n/40 \rfloor + 1}^n z_i^2 + f_{\text{opt}}$ | $z = Q A_{10}^{\text{opt}} R(x - x^{\text{opt}})$ with $R = P_{11} B_{12}$, $Q = P_{21} B_{22}$ |
| Different Powers Function    | $f_{14}(x) = \gamma(n) \times \sum_{i=1}^n (z_i^{2+3 \times \frac{1}{n}}) + f_{\text{opt}}$ | $z = R(x - x^{\text{opt}})$ with $R = P_{11} B_{12}$ |
Table 2: Function descriptions of the multi-modal function group with adequate global structure of the \texttt{bbob-largescale} test suite.

| Group 4: Multi-modal functions with adequate global structure | Formulation | Transformations |
|---------------------------------------------------------------|-------------|----------------|
| Rastrigin Function                                            | \( f_{15}(x) = \gamma(n) \times (10n - 10 \sum_{i=1}^{n} \cos(2\pi z_i) + ||z||^2) + f_{\text{opt}} \) | \( z = RA^{0.1}Q^{0.1}\times(T_{\text{out}}(R(x - x^{\text{opt}}))) \) with \( R = P_{11}B_1P_{12}, Q = P_{21}B_2P_{22} \) |
| Weierstrass Function                                          | \( f_{16}(x) = 10 \left( \frac{1}{n} \sum_{i=1}^{n-1} \sum_{k=0}^{11} \frac{1}{2^k} \cos(2\pi 3^k(z_i + 1/2)) - f_0 \right)^2 + \frac{10}{n} f_{\text{pen}}(x) + f_{\text{opt}} \) | \( z = RA^{1/100}Q^{1/100}\times(T_{\text{out}}(R(x - x^{\text{opt}}))) \) with \( R = P_{11}B_1P_{12}, Q = P_{21}B_2P_{22}, f_0 = \sum_{k=0}^{11} \frac{1}{2^k} \cos(\pi 3^k) \) |
| Schaffers F7 Function                                         | \( f_{17}(x) = \left( \frac{1}{n-1} \sum_{i=1}^{n-1} \left( \sqrt{s_i} + \sqrt{s_i^2 \sin^2(50(s_i)^{1/5})) \right) \right)^2 + 10f_{\text{pen}}(x) + f_{\text{opt}} \) | \( z = A^{1000}Q^{0.5\times}(R(x - x^{\text{opt}})) \) with \( R = P_{11}B_1P_{12}, Q = P_{21}B_2P_{22}, s_i = \sqrt{s_i^2 + s_{i+1}^2}, i = 1, \ldots, n-1 \) |
| Schaffers F7 Function, moderately ill-conditioned             | \( f_{18}(x) = \left( \frac{1}{n-1} \sum_{i=1}^{n-1} \left( \sqrt{s_i} + \sqrt{s_i^2 \sin^2(50(s_i)^{1/5})) \right) \right)^2 + 10f_{\text{pen}}(x) + f_{\text{opt}} \) | \( z = A^{1000}Q^{0.5\times}(R(x - x^{\text{opt}})) \) with \( R = P_{11}B_1P_{12}, Q = P_{21}B_2P_{22}, s_i = \sqrt{s_i^2 + s_{i+1}^2}, i = 1, \ldots, n-1 \) |
| Composite Griewank-Rosenbrock Function F8F2                   | \( f_{19}(x) = \frac{10}{n-1} \sum_{i=1}^{n-1} \left( \frac{s_i}{4000} - \cos(s_i) \right) + 10 + f_{\text{opt}} \) | \( z = \max(1, \frac{\sqrt{n}}{8}) Rx + \frac{1}{2} \) with \( R = P_{11}BP_2, s_i = 100(z_i^2 - z_{i+1})^2 + (z_{i+1} - 1)^2 \), for \( i = 1, \ldots, n-1, x^{\text{opt}} = 1 \) |
Table 3: Function descriptions of the multi-modal function group with weak global structure of the \texttt{bbob-largestcale} test suite.

| Group 5: Multi-modal functions with weak global structure | Formulation | Transformations |
|----------------------------------------------------------|-------------|----------------|
| Schwefel Function                                         | $f_{20}(x) = -\frac{1}{1000^n} \sum_{i=1}^{n} z_i \sin\left(\sqrt{|z_i|}\right) + 4.189828776 \times 10^{-10}$ | $x = 2 \times 1^+ \otimes x$, $z_1 = \hat{x}_1, \hat{x}_{i+1} = \hat{x}_{i+1} + 0.25 \left(\hat{x}_{i} - 2 \cdot x_{i}^{\text{opt}}\right)$, for $i = 1, ..., n - 1$, $z = 100 \left(10^3 \left(\hat{x}_2 - 2 \cdot x_{i}^{\text{opt}}\right) + 2 \cdot \left|x_{i}^{\text{opt}}\right|\right)$, $x_{i}^{\text{opt}} = 4.2096874633 / 10^+$ |
| Gallagher’s Gaussian 101-me Peaks Function                | $f_{21}(x) = T_{\text{osz}} \left(10 - \max_{i=1}^{101} w_i \exp\left(-\frac{1}{2n^2} (z - y_i)^T B_i^T C_i B_i (z - y_i)\right)\right)^2 + f_{\text{pen}}(x) + f_{\text{opt}}$ | $w_i = \left\{ \begin{array}{ll} 1.1 + 8 \times \frac{i - 2}{99} & \text{for } 2 \leq i \leq 101 \\ 10 & \text{for } i = 1 \end{array} \right.$ $B$ is a block-diagonal matrix without permutations of the variables. $C_i = \Lambda_{\alpha_i} / \alpha_i^{1/4}$ where $\Lambda_{\alpha_i}$ is defined as usual, but with randomly permuted diagonal elements. For $i = 2, \ldots, 101$, $\alpha_i$ is drawn uniformly from the set $\{1000^{2/3}, 22, \ldots, 101\}$ without replacement, and $\alpha_i = 1000$ for $i = 1$. The local optima $y_i$ are uniformly drawn from the domain $[-5, 5]^n$ for $i = 2, \ldots, 101$ and $y_1 \in [-4, 4]^n$. The global optimum is at $x_{i}^{\text{opt}} = y_1$. |
| Gallagher’s Gaussian 21-hi Peaks Function                 | $f_{22}(x) = T_{\text{osz}} \left(10 - \max_{i=1}^{21} w_i \exp\left(-\frac{1}{2n^2} (z - y_i)^T B_i^T C_i B_i (z - y_i)\right)\right)^2 + f_{\text{pen}}(x) + f_{\text{opt}}$ | $w_i = \left\{ \begin{array}{ll} 1.1 + 8 \times \frac{i - 2}{19} & \text{for } 2 \leq i \leq 21 \\ 10 & \text{for } i = 1 \end{array} \right.$ $B$ is a block-diagonal matrix without permutations of the variables. $C_i = \Lambda_{\alpha_i} / \alpha_i^{1/4}$ where $\Lambda_{\alpha_i}$ is defined as usual, but with randomly permuted diagonal elements. For $i = 2, \ldots, 21$, $\alpha_i$ is drawn uniformly from the set $\{1000^{2/3}, 22, \ldots, 101\}$ without replacement, and $\alpha_i = 1000^2$ for $i = 1$. The local optima $y_i$ are uniformly drawn from the domain $[-4.9, 4.9]^n$ for $i = 2, \ldots, 21$ and $y_1 \in [-3.92, 3.92]^n$. The global optimum is at $x_{i}^{\text{opt}} = y_1$. |
| Katsuura Function                                         | $f_{23}(x) = \frac{10}{n^2} \prod_{i=1}^{n} \left(1 + \sum_{j=1}^{32} \frac{|2^j z_i - 2^j z_i^{\text{opt}}|}{2^j}\right)^{10/n+2} - 10 / n^2$ | $z = Q A^{100} R (x - x_{i}^{\text{opt}})$ with $R = P_{11} B_1 P_{12}, Q = P_{21} B_2 P_{22}$ |
| Lunacek bi-Rastrigin Function                             | $f_{24}(x) = \gamma(n) \times \left( \min \left( \sum_{i=1}^{n} (\hat{x}_i - \mu_3)^2, n \cdot \sum_{i=1}^{n} (\hat{x}_i - \mu_3)^2 + 10(n - \sum_{i=1}^{n} \cos(2\pi z_i)) \right) + 10^4 f_{\text{pen}}(x) + f_{\text{opt}} \right)$ | $x = 2 \sin(\pi x_{i}^{\text{opt}}) \otimes x$, $x_{i}^{\text{opt}} = 0.5 \mu_0 1^+$, $z = Q A^{100} R (x - \mu_0 1)$ with $R = P_{11} B_1 P_{12}, Q = P_{21} B_2 P_{22}, \mu_0 = 2.5, \mu_1 = -\sqrt{\frac{\mu_0^2 - 1}{s}}, s = 1 - \frac{1}{2\sqrt{n + 20} - 8.2}$ |