CLIQUE MINORS IN CARTESIAN PRODUCTS OF GRAPHS

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ABSTRACT. A clique minor in a graph $G$ can be thought of as a set of connected subgraphs in $G$ that are pairwise disjoint and pairwise adjacent. The Hadwiger number $\eta(G)$ is the maximum cardinality of a clique minor in $G$. It is one of the principle measures of the structural complexity of a graph.

This paper studies clique minors in the Cartesian product $G \square H$. Our main result is a rough structural characterisation theorem for Cartesian products with bounded Hadwiger number. It implies that if the product of two sufficiently large graphs has bounded Hadwiger number then it is one of the following graphs:

• a planar grid with a vortex of bounded width in the outerface,
• a cylindrical grid with a vortex of bounded width in each of the two ‘big’ faces, or
• a toroidal grid.

Motivation for studying the Hadwiger number of a graph includes Hadwiger’s Conjecture, which states that the chromatic number $\chi(G) \leq \eta(G)$. It is open whether Hadwiger’s Conjecture holds for every Cartesian product. We prove that $G \square H$ (where $\chi(G) \geq \chi(H)$) satisfies Hadwiger’s Conjecture whenever:

• $H$ has at least $\chi(G) + 1$ vertices, or
• the treewidth of $G$ is sufficiently large compared to $\chi(G)$.

On the other hand, we prove that Hadwiger’s Conjecture holds for all Cartesian products if and only if it holds for all $G \square K_2$. We then show that $\eta(G \square K_2)$ is tied to the treewidth of $G$.

We also develop connections with pseudoachromatic colourings and connected dominating sets that imply near-tight bounds on the Hadwiger number of grid graphs (Cartesian products of paths) and Hamming graphs (Cartesian products of cliques).

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1. INTRODUCTION

A clique minor in a graph $G$ can be thought of as a set of connected subgraphs in $G$ that are pairwise disjoint and pairwise adjacent. The Hadwiger number $\eta(G)$ is the maximum cardinality of a clique minor in $G$. It is one of the principle measures of the structural complexity of a graph.

Robertson and Seymour [41] proved a rough structural characterisation of graphs with bounded Hadwiger number. It says that such a graph can be constructed by a combination of four ingredients: graphs embedded in a surface of bounded genus, vortices of bounded width inside a face, the addition of a bounded number of apex vertices, and the clique-sum operation. Moreover, each of these ingredients is essential. This result is at the heart of Robertson and Seymour’s proof of Wagner’s Conjecture [42]: Every infinite set of finite graphs contains two graphs, one of which is a minor of the other.

This paper studies clique minors in the (Cartesian) product $G \square H$. Our main result is a rough structural characterisation of products with bounded Hadwiger number, which is less rough than the far more general result by Robertson and Seymour. It says that for connected graphs $G$ and $H$, each with at least one edge, $G \square H$ has bounded Hadwiger number if and only if at least one of the following conditions are satisfied:

- $G$ has bounded treewidth and $H$ has bounded order,
- $H$ has bounded treewidth and $G$ has bounded order, or
- $G$ has bounded hangover and $H$ has bounded hangover,

where hangover is a parameter defined in Section 11. Basically, a graph with bounded hangover is either a cycle or consists of a path of degree-2 vertices joining two connected subgraphs of bounded order with no edge between the subgraphs. This implies that if the product of two sufficiently large graphs has bounded Hadwiger number then it is one of the following graphs:

- a planar grid (the product of two paths) with a vortex of bounded width in the outerface,
- a cylindrical grid (the product of a path and a cycle) with a vortex of bounded width in each of the two ‘big’ faces, or
- a toroidal grid (the product of two cycles).

The key case for the proof of this structure theorem is when $G$ and $H$ are trees. This case is handled in Section 9. The proof for general graphs is given in Sections 10 and 11.

Before proving our main results we develop connections with pseudoachromatic colourings (Section 4) and connected dominating sets (Sections 5 and 6) that imply near-tight bounds on the Hadwiger number of grid graphs (products of paths; Sections 3, 4 and 6) and Hamming graphs (products of cliques; Section 7). As summarised in Table 1, in each case, we improve the best previously known lower bound by a factor of between $\Omega(n^{1/2})$ and $\Omega(n^{3/2})$ to conclude asymptotically tight bounds for fixed $d$.

1.1. Hadwiger’s Conjecture. Motivation for studying clique minors includes Hadwiger’s Conjecture, a far reaching generalisation of the 4-colour theorem, which states that the
chromatic number $\chi(G) \leq \eta(G)$ for every graph $G$. It is open whether Hadwiger’s Conjecture holds for every product. The following classes of products are known to satisfy Hadwiger’s Conjecture (where $G$ and $H$ are connected and $\chi(G) \geq \chi(H)$):

- The product of sufficiently many graphs relative to their maximum chromatic number satisfies Hadwiger’s Conjecture [7].
- If $\chi(H)$ is not too small relative to $\chi(G)$, then $G \square H$ satisfies Hadwiger’s Conjecture [5, 38].

See Section 12 for precise versions of this statements. We add to this list as follows:

- If $H$ has at least $\chi(G) + 1$ vertices, then $G \square H$ satisfies Hadwiger’s Conjecture (Theorem 12.4).
- If the treewidth of $G$ is sufficiently large compared to $\chi(G)$, then $G \square H$ satisfies Hadwiger’s Conjecture (Theorem 12.3).

On the other hand, we prove that Hadwiger’s Conjecture holds for all $G \square K_2$ with $\chi(G) \geq \chi(H)$ if and only if Hadwiger’s Conjecture holds for $G \square K_2$. We then show that $\eta(G \square K_2)$ is tied to the treewidth of $G$. All these results are presented in Section 12.

Clique minors in products have been previously considered by a number of authors [1, 3, 7, 26, 32, 33, 38, 55]. In related work, Xu and Yang [53] and Špacapan [48] studied the connectivity of products, and Drier and Linial [14] studied minors in lifts of graphs.

2. Preliminaries

All graphs considered in this paper are undirected, simple, and finite; see [2, 11]. Let $G$ be a graph with vertex $V(G)$ and edge set $E(G)$. Let $v(G) = |V(G)|$ and $e(G) = |E(G)|$ respectively denote the order and size of $G$. Let $\Delta(G)$ denote the maximum degree of $G$.

Let $K_n$ be the complete graph with $n$ vertices. A clique of a graph $G$ is a complete subgraph of $G$. The clique number of $G$, denoted by $\omega(G)$, is the maximum order of a clique of $G$. Let $P_n$ be the path with $n$ vertices. By default, $V(K_n) = [n]$ and $P_n = (1, 2, \ldots, n)$. A leaf in a graph is a vertex of degree 1. Let $S_n$ be the star graph with $n$ leaves; that is, $S_n = K_{1,n}$.

Consider a graph $H$ embedded in a surface; see [36]. Let $(v_1, v_2, \ldots, v_k)$ be a facial cycle in $H$. Consider a graph $G$ obtained from $H$ by adding sets of vertices $S_1, S_2, \ldots, S_k$ (called bags), such that for each $i \in [k]$ we have $v_i \in S_i \cap V(H) \subseteq \{v_1, \ldots, v_k\}$, and for each vertex $v \in \cup_i S_i$, if $R(v) := \{i \in [k], v \in S_i\}$ then for some $i, j$, either $R(v) = [i, j]$ or

| graph | $d$ | previous best | new result | reference |
|-------|-----|---------------|------------|-----------|
| grid graph $P_n^d$ | even | $\Omega(n^{(d-2)/2})$ | $\Theta(n^{d/2})$ | Theorem 3.2 |
| grid graph $P_n^d$ | odd | $\Omega(n^{(d-1)/2})$ | $\Theta(n^{d/2})$ | Theorem 4.4 |
| Hamming graph $K_n^d$ | even | $\Omega(n^{(d-2)/2})$ | $\Theta(n^{(d+1)/2})$ | Theorem 7.5 |
| Hamming graph $K_n^d$ | odd | $\Omega(n^{(d-1)/2})$ | $\Theta(n^{(d+1)/2})$ | Theorem 7.5 |
Lemma 2.1. If $H$ is a minor of a connected graph $G$, then $G$ has an $H$-minor such that every vertex of $G$ is in some branch set.

Proof. Start with an $H$-minor of $G$. If some vertex of $G$ is not in a branch set, then since $G$ is connected, some vertex $v$ of $G$ is not in a branch set and is adjacent to a vertex that
is in a branch set \(X\). Adding \(v\) to \(X\) gives an \(H\)-minor using more vertices of \(G\). Repeat until every vertex of \(G\) is in some branch set. \(\square\)

To prove the tightness of our lower bound constructions, we use the following elementary upper bounds on the Hadwiger number.

**Lemma 2.2.** For every connected graph \(G\),

\[
\begin{align*}
(a) \quad \eta(G) & \leq \sqrt{2e(G) - 2v(G) + \frac{9}{4}} + \frac{3}{2} \\
(b) \quad \eta(G) & < \sqrt{2e(G) + 1} \\
(c) \quad \eta(G) & \leq \sqrt{\Delta(G) - 2}v(G) + 3.
\end{align*}
\]

**Proof.** Let \(k := \eta(G)\). Say \(X_1, \ldots, X_k\) are the branch sets of \(K_k\)-minor in \(G\). By Lemma 2.1, we may assume that every vertex is in some branch set. Since at least \(\binom{k}{2}\) edges have endpoints in distinct branch sets,

\[e(G) \geq \binom{k}{2} + \sum_{i=1}^{k} e(X_i) \geq \binom{k}{2} + \sum_{i=1}^{k} (v(X_i) - 1) = \binom{k}{2} + v(G) - k.
\]

Thus \(k^2 - 3k - 2e(G) + 2v(G) \leq 0\), and (a) follows by the quadratic formula. Parts (b) and (c) are straightforward consequences of (a). \(\square\)

Note that Lemma 2.2(a) is tight for every tree \((\eta = 2)\). The following result, first proved by Ivančo [26], is another elementary upper bound on \(\eta(G)\). It is tight for a surprisingly large class of graphs; see Proposition 8.3. We include the proof for completeness.

**Lemma 2.3 ([26, 49]).** For every graph \(G\),

\[\eta(G) \leq \left\lfloor \frac{v(G) + \omega(G)}{2} \right\rfloor.
\]

**Proof.** Consider a \(K_n\)-minor in \(G\), where \(n := \eta(G)\). For \(j \geq 1\), let \(n_j\) be the number of branch sets that contain exactly \(j\) vertices. Thus

\[v(G) - n_1 \geq \sum_{j \geq 2} j \cdot n_j \geq 2 \sum_{j \geq 2} n_j = 2(n - n_1).
\]

Hence \(v(G) + n_1 \geq 2n\). The branch sets that contain exactly one vertex form a clique. Thus \(n_1 \leq \omega(G)\) and \(v(G) + \omega(G) \geq 2n\). The result follows. \(\square\)

### 3. Hadwiger Number of Grid Graphs

In this section we consider the Hadwiger number of the products of paths, so called grid graphs. First consider the \(n \times m\) grid \(P_n \square P_m\). It has no \(K_5\)-minor since it is planar. In fact, \(\eta(P_n \square P_m) = 4\) for all \(n \geq m \geq 3\). Similarly, \(P_n \square P_2\) has no \(K_4\)-minor since it is outerplanar, and \(\eta(P_n \square P_2) = 3\) for all \(n \geq 2\).

Now consider the double-grid \(P_n \square P_m \square P_2\), where \(n \geq m \geq 2\). For \(i \in [n]\), let \(C_i\) be the \(i\)-th column in the base copy of \(P_n \square P_m\); that is, \(C_i := \{(i, y, 1) : y \in [m]\}\). For \(j \in [m]\),
let \( R_j \) be the \( j \)-th row in the top copy of \( P_n \square P_m \); that is, \( R_j := \{(x, j, 2) : x \in [n]\} \). Since each \( R_i \) and each \( C_j \) are adjacent, contracting each \( R_i \) and each \( C_j \) gives a \( K_{n,m} \)-minor. Chandran and Sivadasan [7] studied the case \( n = m \), and observed that a \( K_m \)-minor is obtained by contracting a matching of \( m \) edges in \( K_{m,m} \). In fact, contracting a matching of \( m - 1 \) edges in \( K_{n,m} \) gives a \( K_{m+1} \)-minor. (In fact, \( \eta(K_{m,m}) = m + 1 \); see [52] for example.) Now observe that \( R_1 \) is adjacent to \( R_2 \) and \( C_1 \) is adjacent to \( C_2 \). Thus contracting each edge of the matching \( R_3C_3, R_4C_4, \ldots, R_mC_m \) gives a \( K_{m+2} \)-minor, as illustrated in Figure 1.

**Figure 1.** A \( K_{m+2} \)-minor in \( P_m \square P_m \square P_2 \).

Theorem 10.2 below implies an upper bound of \( \eta(P_n \square P_m \square P_2) \leq 2m + 2 \) regardless of \( n \). In fact, Theorem 10.2 implies that \( \eta(P_n \square K_m \square P_2) \leq 2m + 2 \). Summarising,

\[
1 \leq \eta(P_n \square P_m \square P_2) \leq 2m + 2.
\]

We conjecture that the lower bound in (1) is the answer; that is, \( \eta(P_n \square P_m \square P_2) = m+2 \).

The above construction of a clique minor in the double-grid generalises as follows.

**Proposition 3.1.** For all connected graphs \( G \) and \( H \), each with at least one edge,

\[
\eta(G \square H \square P_2) \geq \omega(G) + \omega(H) + \min\{v(G) - \omega(G), v(H) - \omega(H)\}
\]

\[
\geq \min\{v(G), v(H)\} + \min\{\omega(G), \omega(H)\}
\]

\[
\geq \min\{v(G), v(H)\} + 2.
\]

**Proof.** Let \( P \) be a maximum clique of \( G \). Let \( Q \) be a maximum clique of \( H \). Without loss of generality, \( n := v(G) - \omega(G) \leq v(H) - \omega(H) \). Say \( V(G) - V(P) = \{v_1, v_2, \ldots, v_n\} \) and \( V(H) - V(Q) = \{w_1, w_2, \ldots, w_m\} \), where \( n \leq m \). Let \( V(P_2) = \{1, 2\} \).

\footnote{In the case \( n = m \), Chandran and Sivadasan [7] claimed an upper bound of \( \eta(P_m \square P_m \square P_2) \leq 2m + 2 \) without proof.}
For \( x \in V(G) \), let \( A(x) \) be the subgraph of \( G \square H \square P_2 \) induced by \( \{(x, y, 1) : y \in V(H)\} \). For \( y \in V(G) \), let \( B(y) \) be the subgraph of \( G \square H \square P_2 \) induced by \( \{(x, y, 2) : x \in V(G)\} \). Note that each subgraph \( A(x) \) is isomorphic to \( H \), and is thus connected. Similarly, each subgraph \( B(y) \) is isomorphic to \( G \), and is thus connected.

Distinct subgraphs \( A(x) \) and \( A(x') \) are disjoint since the first coordinate of every vertex in \( A(x) \) is \( x \). Distinct subgraphs \( B(y) \) and \( B(y') \) are disjoint since the second coordinate of every vertex in \( B(y) \) is \( y \). Subgraphs \( A(x) \) and \( B(y) \) are disjoint since the third coordinate of every vertex in \( A(x) \) is 1, and the third coordinate of every vertex in \( B(y) \) is 2.

Since the vertex \( (x, y, 1) \) in \( A(x) \) is adjacent to the vertex \( (x, y, 1) \) in \( B(y) \), the \( A(x) \) and \( B(y) \) subgraphs are the branch sets of a complete bipartite \( K_{\omega(G),\omega(H)} \)-minor in \( G \square H \). Moreover, for distinct vertices \( x \) and \( x' \) in the clique \( P \), for any vertex \( y \in V(H) \), the vertex \( (x, y, 1) \) in \( A(x) \) is adjacent to the vertex \( (x', y, 1) \) in \( A(x') \). Similarly, for distinct vertices \( y \) and \( y' \) in the clique \( Q \), for any vertex \( x \in V(G) \), the vertex \( (x, y, 2) \) in \( B(y) \) is adjacent to the vertex \( (x, y', 2) \) in \( B(y') \). For each \( i \in [n] \), let \( X \) be the subgraph induced by \( A(v_i) \cup B(w_i) \). Now \( X(i) \) is connected, since the vertex \( (v_i, w_i, 1) \) in \( A(i) \) is adjacent to the vertex \( (v_i, w_i, 2) \) in \( B(i) \).

We have shown that \( \{A(x) : x \in P\} \cup \{B(x) : x \in Q\} \cup \{X_i : i \in [n]\} \) is a set of \( \omega(G) + \omega(H) + n \) connected subgraphs, each pair of which are disjoint and adjacent. Hence these subgraphs are the branch sets of a clique minor in \( G \square H \). Therefore \( \eta(G \square H \square P_2) \geq \omega(G) + \omega(H) + n \), as desired. The final claims are easily verified.

Now consider the Hadwiger number of the \( d \)-dimensional grid graph

\[
P_n^d := P_n \square P_n \square \cdots \square P_n.
\]

The best previously known bounds are due to Chandran and Sivadasan [7] who proved that

\[
n^{\lfloor (d-1)/2 \rfloor} \leq \eta(P_n^d) \leq \sqrt{2d} n^{d/2} + 1.
\]

In the case that \( d \) is even we now improve this lower bound by a \( \Theta(n) \) factor, and thus determine \( \eta(P_n^d) \) to within a factor of \( 4\sqrt{2d} \) (ignoring lower order terms).

**Theorem 3.2.** For every integer \( n \geq 2 \) and even integer \( d \geq 4 \),

\[
\frac{1}{4} n^{d/2} - O(n^{d/2-1}) \leq \eta(P_n^d) < \sqrt{2d - 2} n^{d/2} + 3.
\]

**Proof.** The upper bound follows from Lemma 2.2(1) since \( \nu(P_n^d) = n^d \) and \( \Delta(P_n^d) = 2d \).

Now we prove the lower bound. Let \( V(P_n^d) = [n]^d \), where two vertices are adjacent if and only if they share \( d - 1 \) coordinates in common and differ by 1 in the remaining coordinate. Let \( p := \frac{d}{4} \).

For \( j_1, j_2 \in \left[ \frac{n}{2} \right] \) and \( j_3, \ldots, j_p \in [n] \), let \( A(j_1, \ldots, j_p) \) be the subgraph of \( P_n^d \) induced by

\[
\{(2j_1, 2j_2, j_3, \ldots, j_p, x_1, x_2, \ldots, x_p) : x_i \in [n], i \in [p]\}.
\]
let $B\langle j_1, \ldots, j_p \rangle$ be the subgraph of $P_n^d$ induced by
\[
\{(2x_1 - 1, x_2, x_3, \ldots, x_p, j_1, j_2, \ldots, j_p) : x_1 \in [\frac{n}{4}], x_i \in [n], i \in [2, p] \}
\]
\[
\cup \{(x_1, 2x_2 - 1, x_3, x_4, \ldots, x_p, j_1, j_2, \ldots, j_p) : x_1 \in [n], x_2 \in [\frac{n}{4}], x_i \in [n], i \in [3, p] \};
\]
and let $X\langle j_1, \ldots, j_p \rangle$ be the subgraph of $P_n^d$ induced by $A\langle j_1, \ldots, j_p \rangle \cup B\langle j_1, \ldots, j_p \rangle$.

Each $A\langle j_1, \ldots, j_p \rangle$ subgraph is disjoint from each $B\langle j'_1, \ldots, j'_p \rangle$ subgraph since the first and second coordinates of each vertex in $A\langle j_1, \ldots, j_p \rangle$ are both even, while the first or second coordinate of each vertex in $B\langle j'_1, \ldots, j'_p \rangle$ is odd. Two distinct subgraphs $A\langle j_1, \ldots, j_p \rangle$ and $A\langle j'_1, \ldots, j'_p \rangle$ are disjoint since the $p$-tuples determined by the first $p$ coordinates are distinct. Similarly, two distinct subgraphs $B\langle j_1, \ldots, j_p \rangle$ and $B\langle j'_1, \ldots, j'_p \rangle$ are disjoint since the $p$-tuples determined by the last $p$ coordinates are distinct. Hence each pair of distinct subgraphs $X\langle j_1, \ldots, j_p \rangle$ and $X\langle j'_1, \ldots, j'_p \rangle$ are disjoint.

Observe that $A\langle j_1, \ldots, j_p \rangle$ is isomorphic to $P_n^p$, and is thus connected. In particular, every pair of vertices
\[
(2j_1, 2j_2, j_3, \ldots, j_p, x_1, x_2, \ldots, x_p)
\]
and
\[
(2j_1, 2j_2, j_3, \ldots, j_p, x'_1, x'_2, \ldots, x'_p)
\]
in $A\langle j_1, \ldots, j_p \rangle$ are connected by a path of length $\sum_i |x_i - x'_i|$ contained in $A\langle j_1, \ldots, j_p \rangle$.

To prove that $B\langle j_1, \ldots, j_p \rangle$ is connected, consider a pair of vertices
\[
v = (x_1, x_2, \ldots, x_p, j_1, j_2, \ldots, j_p)
\]
and
\[
v' = (x'_1, x'_2, \ldots, x'_p, j_1, j_2, \ldots, j_p)
\]
in $B\langle j_1, \ldots, j_p \rangle$. If $x_1$ is even then walk along any one of the dimension-1 edges incident to $v$. This neighbour is in $B\langle j_1, \ldots, j_p \rangle$, and its first coordinate is odd. Thus we can now assume that $x_1$ is odd. Similarly, we can assume that $x_2$, $x'_1$, and $x'_2$ are all odd. Then
\[
(x_1, x_2, \ldots, x_p, j_1, j_2, \ldots, j_p)
\]
and
\[
(x'_1, x'_2, \ldots, x'_p, j_1, j_2, \ldots, j_p)
\]
are connected by a path of length $\sum_i |x_i - x'_i|$ contained in $B\langle j_1, \ldots, j_p \rangle$. Thus $B\langle j_1, \ldots, j_p \rangle$ is connected. The vertex
\[
(2j_1, 2j_2, j_3, \ldots, j_p, j_1, j_2, \ldots, j_p)
\]
in $A\langle j_1, \ldots, j_p \rangle$ is adjacent to the vertex
\[
(2j_1 - 1, 2j_2, j_3, \ldots, j_p, j_1, j_2, \ldots, j_p)
\]
in $B\langle j_1, \ldots, j_p \rangle$. Thus $X\langle j_1, \ldots, j_p \rangle$ is connected. Each pair of subgraphs $X\langle j_1, \ldots, j_p \rangle$ and $X\langle j'_1, \ldots, j'_p \rangle$ are adjacent since the vertex
\[
(2j_1, 2j_2, j_3, \ldots, j_p, j'_1, j'_2, \ldots, j'_p)
\]
in $A\langle j_1, \ldots, j_p \rangle$ is adjacent to the vertex
\[
(2j_1 - 1, 2j_2, j_3, \ldots, j_p, j'_1, j'_2, \ldots, j'_p)
\]
in $B\langle j'_1, \ldots, j'_p \rangle$.

Hence the $X\langle j_1, \ldots, j_p \rangle$ form a complete graph minor in $P_n^d$ of order $n^{d/2-1}|\frac{n}{4}|^2 = \frac{1}{16} n^{d/2} - O(n^{d/2-1})$.  \(\square\)
Note that for particular values of \( n \), the lower bound in Theorem 3.2 is improved by a constant factor in Corollary 6.6 below.

4. Odd-Dimensional Grids and Pseudoachromatic Colourings

The ‘dimension pairing’ technique used in Section 3 to construct large clique minors in even-dimensional grids does not give tight bounds for odd-dimensional grids. To construct large clique minors in odd-dimensional grids we use the following idea.

A pseudoachromatic \( k \)-colouring of a graph \( G \) is a function \( f : V(G) \to [k] \) such that for all distinct \( i, j \in [k] \) there is an edge \( uv \in E(G) \) with \( f(v) = i \) and \( f(w) = j \).

The pseudoachromatic number of \( G \), denoted by \( \psi(G) \), is the maximum integer \( k \) such that there is a pseudoachromatic \( k \)-colouring of \( G \). Pseudoachromatic colourings were introduced by Gupta [19] in 1969, and have since been widely studied. For example, many authors [17, 25, 34, 54] have proved\(^2\) that

\[
\psi(P_n) > \sqrt{2n - 2} - 2. \tag{2}
\]

Note that the only difference between a pseudoachromatic colouring and a clique minor is that each colour class is not necessarily connected. We now show that the colour classes in a pseudoachromatic colouring can be made connected in a three-dimensional product.

**Theorem 4.1.** Let \( G, H, \) and \( I \) be connected graphs. Let \( A \) be a minor of \( G \) with at least two vertices in each branch set. Let \( B \) be a minor of \( H \) with at least two vertices in each branch set. Let \( C \) be a minor of \( I \) with at least two vertices in each branch set. If \( v(B) \geq v(C) \) then

\[
\eta(G \square H \square I) \geq \min\{\psi(A), v(B)\} \cdot v(C). \tag{3}
\]

**Proof.** (The reader should keep the example of \( G = H = I = P_n \) and \( A = B = C = P_{[n/2]} \) in mind.)

Let \( V(B) = \{y_1, \ldots, y_{v(B)}\} \) and \( V(C) = \{z_1, \ldots, z_{v(C)}\} \). By contracting edges in \( H \), we may assume that there are exactly two vertices of \( H \) in each branch set of \( B \). Label the two vertices of \( H \) in the branch set corresponding to each \( y_j \) by \( y_j^+ \) and \( y_j^- \). By contracting edges in \( I \), we may assume that there are exactly two vertices of \( I \) in each branch set of \( C \). Label the two vertices of \( I \) in the branch set corresponding to each \( z_j \) by \( z_j^+ \) and \( z_j^- \).

Let \( k := \min\{\psi(A), v(B)\} \). Let \( f : V(A) \to [k] \) be a pseudoachromatic colouring of \( A \). Our goal is to prove that \( \eta(G \square H \square I) \geq k \cdot v(C) \).

By contracting edges in \( G \), we may assume that there are exactly two vertices of \( G \) in each branch set of \( A \). Now label the two vertices of \( G \) in the branch set corresponding

\(^2\)For completeness, we prove that \( \psi(P_n) > \sqrt{2n - 2} - 2 \). Let \( P_n = (x_1, \ldots, x_n) \). Let \( t \) be the maximum odd integer such that \( \binom{n}{t} \leq n - 1 \). Then \( t > \sqrt{2n - 2} - 2 \). Denote \( V(K_t) \) by \( \{v_1, \ldots, v_t\} \). Since \( t \) is odd, \( K_t \) is Eulerian. Orient the edges of \( K_t \) by following an Eulerian cycle \( C = (e_1, e_2, \ldots, e_{\binom{n}{t}}) \). For \( \ell \in [\binom{n}{t}] \), let \( f(x_\ell) = i \), where \( e_\ell = (v_i, v_j) \). For \( \ell \in [\binom{n}{t} + 1, n] \), let \( f(x_\ell) = 1 \). Consider distinct colours \( i, j \in [k] \). Thus for some edge \( e_{\ell+1} \) of \( K_t \), without loss of generality, \( e_{\ell+1} = (v_i, v_j) \). Say \( e_{\ell+1} = (v_j, v_k) \) is the next edge in \( C \), where \( e_{\ell+1} \) means \( e_1 \) if \( \ell = \binom{n}{t} \). Since \( \ell \leq \binom{n}{t} \leq n - 1 \), we have \( \ell + 1 \in [n] \). By construction, \( f(x_\ell) = i \) and \( f(x_{\ell+1}) = j \). Thus \( f \) is a pseudoachromatic colouring.
to each vertex $v$ of $A$ by $v^+$ and $v^-$ as follows. Let $T$ be a spanning tree of $A$. Orient the edges of $T$ away from some root vertex $r$. Arbitrarily label the vertices $r^+$ and $r^-$ of $G$. Let $w$ be a non-root leaf of $T$. Label each vertex of $T - w$ by induction. Now $w$ has one incoming arc $(v, w)$. Some vertex in the branch set corresponding to $v$ is adjacent to some vertex in the branch set corresponding to $w$. Label $w^+$ and $w^-$ so that there is an edge in $G$ between $v^+$ and $w^-$, or between $v^-$ and $w^+$.

For $v \in V(A)$ and $j \in [v(C)]$, let $P(v, j)$ be the subgraph of $G \Box H \Box I$ induced by

$$\{(v^+, y_j^+, z) : z \in V(I)\},$$

and let $Q(v, j)$ be the subgraph of $G \Box H \Box I$ induced by

$$\{(v^-, y, z_j^+) : y \in V(H)\}.$$

For $i \in [k]$ (and thus $i \leq v(B)$) and $j \in [v(C)]$, let $R(i, j)$ be the subgraph of $G \Box H \Box I$ induced by

$$\{(v, y_i^-, z_j^-) : v \in V(G)\},$$

and let $X(i, j)$ be the subgraph of $G \Box H \Box I$ induced by

$$\cup\{P(v, j) \cup Q(v, j) \cup R(i, j) : v \in f^{-1}(i)\},$$

as illustrated in Figure 2. We now prove that the $X(i, j)$ are the branch sets of a clique minor in $G \Box H \Box I$.

![Figure 2. The branch set $X(i, j)$ in Proposition 4.2, where $f^{-1}(i) = \{u, v, w\}$.](image)

First we prove that each $X(i, j)$ is connected. Observe that each $P(v, j)$ is a copy of $I$ and each $Q(i, j)$ is a copy of $H$, and are thus connected. Moreover, the vertex $(v^+, y_j^+, z_j^+)$ in $P(v, j)$ is adjacent to the vertex $(v^-, y_j^+, z_j^+)$ in $Q(v, j)$. Thus $P(v, j) \cup Q(v, j)$ is connected. Now each $R(i, j)$ is a copy of $G$, and is thus connected. For each $v \in f^{-1}(i)$, the vertex $(v^-, y_i^-, z_j^+)$ which is in $Q(v, j)$, is adjacent to $(v^-, y_i^-, z_j^-)$ which is in $R(i, j)$. Thus $X(i, j)$ is connected.

Now consider distinct subgraphs $X(i, j)$ and $X(i', j')$. We first prove that $X(i, j)$ and $X(i', j')$ are disjoint. Distinct subgraphs $P(v, j)$ and $P(w, j')$ are disjoint since the first
two coordinates of every vertex in \( P(v, j) \) are \((v^+, y_j^+)\), which are unique to \((v, j)\). Similarly, distinct subgraphs \( Q(v, j) \) and \( Q(w, j') \) are disjoint since the first and third coordinates of every vertex in \( Q(v, j) \) are \((y_j^-, z_j^+)\), which are unique to \((v, j)\). Every \( P(v, j) \) is disjoint from every \( Q(w, j') \) since the first coordinate of every vertex in \( P(v, j) \) is positive, while the first coordinate of every vertex in \( Q(w, j') \) is negative. Observe that \( R(i, j) \) and \( R(i', j') \) are disjoint since the second and third coordinates of every vertex in \( R(i, j) \) are \((y_i^-, z_j^-)\), which are unique to \((i, j)\). Also \( R(i, j) \) is disjoint from every \( P(v, j') \cup Q(v, j') \) since the second and third coordinate of every vertex in \( R(i, j) \) are negative, while every vertex in \( P(v, j') \cup Q(v, j') \) has a positive second or third coordinate. Therefore distinct \( X(i, j) \) and \( X(i', j') \) are disjoint.

It remains to prove that distinct subgraphs \( X(i, j) \) and \( X(i', j') \) are adjacent. If \( i = i' \) then the vertex \((v^+, y_j^+, z_j^+)\), which is in \( P(i, j) \), is adjacent to the vertex \((v^-, y_j^+, z_j^+)\), which is in \( Q(i', j') \). Now assume that \( i \neq i' \). Then \( f(v) = i \) and \( f(w) = i' \) for some edge \( vw \) of \( A \). By the labelling of vertices in \( G \), without loss of generality, there is an edge in \( G \) between \( v^+ \) and \( w^- \). Thus the vertex \((v^+, y_j^+, z_j^+)\), which is in \( P(v, j) \subset X(i, j) \), is adjacent to the vertex \((w^-, y_j^+, z_j^+)\), which is in \( Q(w, j) \subset X(i', j') \). In both cases, \( X(i, j) \) and \( X(i', j') \) are adjacent.

Hence the \( X(i, j) \) are the branch sets of a clique minor in \( G \oplus H \oplus I \). Thus \( \eta(G \oplus H \oplus I) \geq k \cdot \psi(C) \). \( \square \)

Now consider the Hadwiger number of the three-dimensional grid \( P_n \oplus P_n \oplus P_n \). Prior to this work the best lower and upper bounds on \( \eta(P_n \oplus P_n \oplus P_n) \) were \( \Omega(n) \) and \( O(n^{3/2}) \) respectively \([3, 7, 38]\). The next result improves this lower bound by a \( \Theta(\sqrt{n}) \) factor, thus determining \( \eta(P_n \oplus P_n \oplus P_n) \) to within a factor of 4 (ignoring lower order terms).

**Proposition 4.2.** For all integers \( n \geq m \geq 1 \),

\[
\frac{1}{2}n\sqrt{m} - O(n + \sqrt{m}) < \eta(P_n \oplus P_n \oplus P_m) \leq 2n\sqrt{m} + 3.
\]

**Proof.** The upper bound follows from Lemma 2.2(c) since \( P_n \oplus P_n \oplus P_m \) has \( n^2m \) vertices and maximum degree 6.

Now we prove the lower bound. \( P_m \) has a \( P_{\lfloor m/2 \rfloor} \)-minor with two vertices in each branch set. By \([2]\), \( \psi(P_{\lfloor m/2 \rfloor}) > \sqrt{m - 3} - 2 \). By Theorem 4.1 with \( G = P_m \), \( A = P_{\lfloor m/2 \rfloor} \), \( H = I = P_n \), and \( B = C = P_{\lfloor n/2 \rfloor} \) (and since \( \psi(P_{\lfloor m/2 \rfloor}) \leq \lfloor m/2 \rfloor \leq \lfloor n/2 \rfloor = \psi(B) \)),

\[
\eta(P_n \oplus P_n \oplus P_m) \geq (\sqrt{m - 3} - 2)\lfloor n/2 \rfloor = \frac{1}{2}n\sqrt{m} - O(n + \sqrt{m})
\]

as desired. \( \square \)

Here is another scenario when tight bounds for three-dimensional grids can be obtained.

**Proposition 4.3.** For all integers \( n \geq m \geq 1 \) such that \( n \leq \frac{1}{4}m^2 \),

\[
\frac{1}{2}m\sqrt{n} - O(m + \sqrt{n}) < \eta(P_n \oplus P_m \oplus P_m) \leq 2m\sqrt{n} + 3.
\]
Proof. The upper bound follows from Lemma 2.2(c) since \( P_n \square P_m \square P_m \) has \( m^2 n \) vertices and maximum degree 6. For the lower bound, apply Theorem 4.1 with \( G = P_n \), \( A = P_{[n/2]} \), \( H = I = P_m \), and \( B = C = P_{[m/2]} \). By (2), \( \psi(P_{[n/2]}) > \sqrt{n - 3} - 2 \). Thus
\[
\eta(P_n \square P_m \square P_m) \geq \min\{(\sqrt{n - 3} - 2), [\frac{m}{2}] \cdot [\frac{m}{2}]\}.
\]
Since \( n \leq \frac{1}{2} m^2 \),
\[
\eta(P_n \square P_m \square P_m) \geq (\sqrt{n - 3} - 2) \cdot [\frac{m}{2}] = \frac{1}{2} m \sqrt{n} - O(m + \sqrt{n})
\]
as desired. \( \square \)

Now consider the Hadwiger number of \( P_n^d \) for odd \( d \). Prior to this work the best lower and upper bounds on \( \eta(P_n^d) \) were \( \Omega(n^{(d-1)/2}) \) and \( O(\sqrt{d} \cdot n^{d/2}) \) respectively [5, 7, 38]. The next result improves this lower bound by a \( \Theta(\sqrt{n}) \) factor, thus determining \( \eta(P_n^d) \) to within a factor of \( 2\sqrt{2(d-1)} \) (ignoring lower order terms).

**Theorem 4.4.** For every integer \( n \geq 1 \) and odd integer \( d \geq 3 \),
\[
\frac{1}{2} n^{d/2} - O(n^{(d-1)/2}) < \eta(P_n^d) \leq \sqrt{2(d-1)} n^{d/2} + 3.
\]

Proof. The upper bound follows from Lemma 2.2(c) since \( \nu(P_n^d) = n^d \) and \( \Delta(P_n^d) = 2d \).

Now we prove the lower bound. Let \( G := P_n \) and \( A := P_{[n/2]} \). By (2), \( \psi(A) > \sqrt{n - 3} - 2 \). Let \( H := P_{[n/2]}^{(d-1)/2} \). Every grid graph with \( m \) vertices has a matching of \( [\frac{m}{2}] \) edges. (Proof: induction on the number of dimensions.) Thus \( H \) has a minor \( B \) with \( [\frac{1}{2} n^{(d-1)/2}] \) vertices, and two vertices in each branch set. By Theorem 4.1 with \( I = H \) and \( C = B \) (and since \( \psi(P_{[n/2]}) \leq [\frac{n}{2}] \leq [\frac{1}{2} n^{(d-1)/2}] \)),
\[
\eta(P_n^d) \geq (\sqrt{n - 3} - 2) [\frac{1}{2} n^{(d-1)/2}] = \frac{1}{2} n^{d/2} - O(n^{(d-1)/2}),
\]
as desired. \( \square \)

### 5. Star Minors and Dominating Sets

Recall that \( S_t \) is the star graph with \( t \) leaves. Consider the Hadwiger number of the product of \( S_t \) with a general graph.

**Lemma 5.1.** For every connected graph \( G \) and for every integer \( t \geq 1 \),
\[
\eta(G \square S_t) \geq \min\{\nu(G), t + 1\}.
\]

Proof. Let \( k := \min\{\nu(G), t + 1\} \). Let \( V(G) := \{v_1, v_2, \ldots, v_n\} \) and \( V(S_t) := \{r\} \cup [t] \), where \( r \) is the root. For \( i \in [k - 1] \), let \( X_i \) be the subgraph of \( G \square S_t \) induced by \( \{(v_i, r)\} \cup \{(v_j, i) : j \in [n]\} \). Since \( G \) is connected, and \((v_i, r)\) is adjacent to \((v_i, i)\), each \( X_i \) is connected. Let \( X_k \) be the subgraph consisting of the vertex \((v_k, r)\). For distinct \( i, j \in [k] \) with \( i < j \) (and thus \( i \in [k - 1] \subseteq [t] \)), the subgraphs \( X_i \) and \( X_j \) are disjoint, and vertex \((v_j, i)\), which is in \( X_i \), is adjacent to \((v_j, i)\), which is in \( X_j \). Thus \( X_1, \ldots, X_k \) are the branch sets of a \( K_k \)-minor in \( G \square S_t \), as illustrated in Figure 3 when \( G \) is a path. \( \square \)
Note that the lower bound in Lemma 5.1 is within a constant factor of the upper bound in Lemma 2.2 whenever \( e(G) = \Theta(v(G)) = \Theta(t) \). In particular, if \( G \) is a tree with at least \( t + 1 \) vertices, then

\[
t + 1 \leq \eta(G \boxtimes S_t) \leq \eta(G \boxtimes K_{t+1}) \leq t + 2,
\]

where the upper bound is proved in Theorem 10.1 below. When \( G \) is another star, Ivančo [26] determined the Hadwiger number precisely. We include the proof for completeness.

**Lemma 5.2** ([26]). For all integers \( n \geq m \geq 2 \),

\[
\eta(S_n \boxtimes S_m) = m + 2.
\]

**Proof.** Let \( V(S_n) := \{r\} \cup [n] \), where \( r \) is the root vertex. Observe that for all \( i \in [n] \) and \( j \in [m] \), vertex \((i, j)\) has degree 2; it is adjacent to \((r, j)\) and \((i, r)\). In every graph except \( K_3 \), contracting an edge incident to a degree-2 vertex does not change the Hadwiger number. Thus replacing the path \((r, j)(i, j)(i, r)\) by the edge \((r, j)(i, r)\) does not change the Hadwiger number. Doing so gives \( K_{1, m,n} \). Ivančo [26] proved that \( \eta(K_{1, m,n}) = m + 2 \). Thus \( \eta(S_n \boxtimes S_m) = m + 2 \). In fact, Ivančo [26] determined the Hadwiger number of every complete multipartite graph (a result rediscovered by the author [52]). \( \square \)

For every graph \( G \), let \( \text{star}(G) \) be the maximum integer \( t \) for which \( S_t \) is minor of \( G \). Since a vertex and its neighbours form a star, \( \text{star}(G) \geq \Delta(G) \). Thus Lemmas 5.1 and 5.2 imply:

**Corollary 5.3.** For all connected graphs \( G \) and \( H \),

\[
\eta(G \boxtimes H) \geq \min\{\text{star}(G) + 1, v(H)\} \geq \min\{\Delta(G) + 1, v(H)\}, \text{ and}
\]

\[
\eta(G \boxtimes H) \geq \min\{\text{star}(G), \text{star}(H)\} + 2 \geq \min\{\Delta(G), \Delta(H)\} + 2.
\]

As an aside, we now show that star minors are related to radius and bandwidth. Let \( G \) be a connected graph. The radius of \( G \), denoted by \( \text{rad}(G) \), is the minimum, taken over all vertices \( v \) of \( G \), of the maximum distance between \( v \) and some other vertex of \( G \).

Each vertex \( v \) that minimises this maximum distance is a centre of \( G \). The bandwidth of
where

\[ X \]

Let \( G \) be a connected graph, and \( W \) alikar [44]). Let \( G \) and \( W \) imply that the product of graphs with small radii or large bandwidth has 

\[ \gamma \]

Laskar [24] proved that when \( G \) is considered to have no leaves, and 

\[ K \]

\[ n \]

\[ b \]

\[ v \]

\[ S \]

\[ i \]

\[ S \]

\[ v \]

\[ S \]

\[|\]

Proof. First we prove (a). Let \( v \) be a centre of \( G \). For \( i \in [0, \text{rad}(G)] \), let \( V_i \) be the set of vertices at distance \( i \) from \( v \). Thus the \( V_i \) are a partition of \( V(G) \), and \(|V_0| = 1\). For each \( i \in [\text{rad}(G)] \), contracting \( V_0, \ldots, V_{i-1} \) into a single vertex and deleting \( V_{i+1}, \ldots, V_{\text{rad}(G)} \) gives a \( S_{|V_i|} \)-minor. Thus \( \text{star}(G) \geq |V_i| \). Hence \( v(G) = \sum_i |V_i| \leq 1 + \text{rad}(G) \cdot \text{star}(G) \).

Now we prove (b). Let \((v_1, \ldots, v_n)\) be a linear ordering of \( V(G) \) such that if \( v_a \in V_i \) and \( v_b \in V_j \) with \( i < j \), then \( a < b \). Consider an edge \( v_a v_b \in E(G) \). Say \( v_a \in V_i \) and \( v_b \in V_j \). Without loss of generality, \( i \leq j \). By construction, \( j \leq i + 1 \). Thus \( b - a \leq |V_i| + |V_{i+1}| - 1 \leq 2 \cdot \text{star}(G) - 1 \). Hence \( \text{bw}(G) \leq 2 \cdot \text{star}(G) - 1. \)

Note that Lemma 5.4(a) is best possible for \( G = P_{2n+1} \), which has \( \text{star}(G) = 2 \) and \( \text{rad}(G) = n \), or for \( G = K_{n+1} \), which has \( \text{star}(G) = n - 1 \) and \( \text{rad}(G) = 1 \). Corollary 5.7 and Lemma 5.4 imply that the product of graphs with small radii or large bandwidth has large Hadwiger number; we omit the details.

Star minors are related to connected dominating sets (first defined by Sampathkumar and Walikar [44]). Let \( G \) be a graph. A set of vertices \( S \subseteq V(G) \) is dominating if each vertex in \( V(G) - S \) is adjacent to a vertex in \( S \). The domination number of \( G \), denoted by \( \gamma(G) \), is the minimum cardinality of a dominating set of \( G \). If \( S \) is dominating and \( G[S] \) is connected, then \( S \) and \( G[S] \) are connected dominating. Only connected graphs have connected dominating sets. The connected domination number of a connected graph \( G \), denoted by \( \gamma_c(G) \), is the minimum cardinality of a connected dominating set of \( G \). Finally, let \( \ell(G) \) be the maximum number of leaves in a spanning tree of \( G \) (where \( K_1 \) is considered to have no leaves, and \( K_2 \) is considered to have one leaf). Hedetniemi and Laskar [24] proved that \( \gamma_c(G) = v(G) - \ell(G) \). We extend this result as follows.

Lemma 5.5. For every connected graph \( G \),

\[ \text{star}(G) = v(G) - \gamma_c(G) = \ell(G). \]

Proof. Let \( T \) be a spanning tree of \( G \) with \( \ell(G) \) leaves. Let \( S \) be the set of non-leaf vertices of \( T \). Thus \( S \) is a connected dominating set of \( G \), implying \( \gamma_c(G) \leq v(G) - \ell(G) \) and \( \ell(G) \leq v(G) - \gamma_c(G) \).

Say \( S \) is a connected dominating set in \( G \) of order \( \gamma_c(G) \). Contracting \( G[S] \) gives a \( S_{v(G)-\gamma_c(G)} \)-minor in \( G \). Thus \( \text{star}(G) \geq v(G) - \gamma_c(G) \).

Now suppose that \( X_0, X_1, \ldots, X_{\text{star}(G)} \) are the branch sets of a \( S_{\text{star}(G)} \)-minor in \( G \), where \( X_0 \) is the root. By Lemma 2.1 we may assume that every vertex of \( G \) is in some \( X_i \). Let \( T_i \) be a spanning tree of each \( G[X_i] \). Each \( T_i \) has at least two leaves, unless \( v(T_i) \leq 2 \). Let \( T \) be the tree obtained from \( \cup_i T_i \) by adding one edge between \( T_0 \) and \( T_i \).
for each \( i \in [\text{star}(G)] \). Each such \( T_i \) contributes at least one leaf to \( T \). Thus \( T \) has at least \( \text{star}(G) \) leaves, implying \( \ell(G) \geq \text{star}(G) \).

**Corollary 5.6.** For every tree \( T \neq K_2 \), \( \text{star}(T) \) equals the number of leaves in \( T \).

Corollary 5.3 and Lemma 5.5 imply:

**Corollary 5.7.** For all connected graphs \( G \) and \( H \),

\[
\eta(G \Box H) \geq \min\{\nu(G) - \gamma_c(G) + 1, \nu(H)\} \quad \text{and} \\
\eta(G \Box H) \geq \min\{\nu(G) - \gamma_c(G), \nu(H) - \gamma_c(H)\} + 2.
\]

We now show that a connected dominating set in a product can be constructed from a dominating set in one of its factors.

**Lemma 5.8.** For all connected graphs \( G \) and \( H \),

\[\gamma_c(G \Box H) \leq (\nu(G) - 1) \cdot \gamma(H) + \nu(H)\]

**Proof.** Let \( S \) be a minimum dominating set in \( H \). Let \( v \) be an arbitrary vertex in \( G \). Consider the set of vertices in \( G \Box H \),

\[T := \{(x, y) : x \in V(G) - \{v\}, y \in S\} \cup \{(v, y) : y \in V(H)\}.
\]

Then \(|T| = (\nu(G) - 1) \cdot \gamma(H) + \nu(H)\). First we prove that \( T \) is dominating. Consider a vertex \((x, y)\) of \( G \Box H \). If \( y \in S \) then \((x, y) \in T\). Otherwise, \( y \) is adjacent to some vertex \( z \in S \), in which case \((x, y)\) is adjacent to \((x, z)\), which is in \( T \). Thus \( T \) is dominating. Now we prove that \( T \) is connected. Since \( G \) is connected, for each vertex \( x \) of \( G \), there is a path \( P(x, v) \) between \( x \) and \( v \) in \( G \). Consider distinct vertices \((x, y)\) and \((x', y')\) in \( G \Box H \). Since \( H \) is connected there is a path \( Q(y, y') \) between \( y \) and \( y' \) in \( H \). Thus

\[\{(s, y) : s \in P(x, v)\} \cup \{(v, t) : t \in Q(y, y')\} \cup \{(s, y') : s \in P(x', v)\}\]

is a path between \((x, y)\) and \((x', y')\) in \( G \Box H \), all of whose vertices are in \( T \). Thus \( T \) is a connected dominating set. \( \Box \)

Corollary 5.7 (with \( G = A \Box B \) and \( H = C \)) and Lemma 5.8 (with \( G = A \) and \( H = B \)) imply:

**Corollary 5.9.** For all connected graphs \( A, B, C \),

\[\eta(A \Box B \Box C) \geq \min\{\nu(A) \cdot \nu(B) - (\nu(A) - 1) \cdot \gamma(B) - \nu(B) + 1, \nu(C)\}.
\]

There are hundreds of theorems about domination in graphs that may be used in conjunction with the above results to construct clique minors in products; see the monograph [23]. Instead, we now apply perhaps the most simple known bound on the order of dominating sets to conclude tight bounds on \( \eta(G^d) \) whenever \( d \geq 4 \) is even and \( e(G) = \Theta(\nu(G)) \).

**Theorem 5.10.** Let \( G \neq K_1 \) be a connected graph with \( n \) vertices and average degree \( \alpha \). Then for every even integer \( d \geq 4 \),

\[
\frac{1}{2} n^{d/2} - n^{d/2-1} + 2 \leq \eta(G^d) \leq \sqrt{\alpha d - 1} n^{d/2} + 3.
\]
Proof. The upper bound follows from Lemma 2.2(a) since
\[ 2e(G^d) = 2d \cdot v(G)^{d-1} \cdot e(G) = d \cdot n^{d-1} \cdot \alpha n = \alpha d \cdot n^d. \]

Now we prove the lower bound. Ore [37] observed that for every connected graph 
\( H \neq K_1 \), the smaller colour class in a 2-colouring of a spanning tree of \( H \) is dominating
in \( H \). Thus \( \gamma(H) \leq \frac{1}{2}v(H) \). This observation with \( H = G^{d/2-1} \) gives
\[ \gamma(G^{d/2-1}) \leq \frac{1}{2}v(G^{d/2-1}) = \frac{1}{2}n^{d/2-1}. \]
By Lemma 5.8
\[ \gamma_c(G^{d/2}) \leq (v(G)-1) \cdot \gamma(G^{d/2-1}) + v(G^{d/2-1}) \leq (n-1)(\frac{1}{2}n^{d/2-1}) + n^{d/2-1} = \frac{1}{2}n^{d/2} + \frac{1}{2}n^{d/2-1}. \]
By Corollary 5.7,
\[ \eta(G^d) \geq v(G^{d/2}) - \gamma_c(G^{d/2}) + 2 \geq n^{d/2} - (\frac{1}{2}n^{d/2} + \frac{1}{2}n^{d/2-1}) = \frac{1}{2}n^{d/2} - \frac{1}{2}n^{d/2-1} + 2, \]
as desired. \( \square \)

6. Dominating Sets and Clique Minors in Even-Dimensional Grids

The results in Section 5 motivate studying dominating sets in grid graphs. First consider
the one-dimensional case of \( P_n \). It is well known and easily proved that \( \gamma(P_n) = \lceil \frac{n}{3} \rceil \). Thus, by Lemma 5.8 for every connected graph \( G \),
\[ \gamma_c(G \sqcap P_n) \leq (v(G) - 1)\lceil \frac{n}{3} \rceil + n. \]
In particular,
\[ \gamma_c(P_m \sqcap P_n) \leq (m - 1)\lceil \frac{n}{3} \rceil + n \leq (m - 1)(\frac{n+2}{3}) + n = \frac{1}{3}(nm + 2m + 2n - 2). \]
Hence Corollary 5.7 with \( G = P_n \sqcap P_m \) implies the following bound on the Hadwiger
number of the 4-dimensional grid:
\[ \eta(P_n \sqcap P_m \sqcap P_n \sqcap P_m) \geq nm - \frac{1}{4}(nm + 2m + 2n - 2) + 2 = \frac{2}{3}(nm - m - n + 4). \]
This result improves upon the bound in Theorem 3.2 by a constant factor.

Dominating sets in 2-dimensional grid graphs are well studied [8, 9, 10, 15, 16, 21, 22, 27, 28, 33, 46, 51]. Using the above technique, these results imply bounds on the
Hadwiger number of the 6-dimensional grid. We omit the details, and jump to the general case.

We first construct a dominating set in a general grid graph.

Lemma 6.1. Fix integers \( d \geq 1 \) and \( n_1, n_2, \ldots, n_d \geq 1 \). Let \( S \) be the set of vertices
\[ S := \{(x_1, x_2, \ldots, x_d) : x_i \in [n_i], i \in [d], \sum_{i \in [d]} i \cdot x_i \equiv 0 \pmod{2d+1}\}. \]
For \( j \in [d] \), let \( B_j \) be the set of vertices
\[ B_j := \{(x_1, \ldots, x_{j-1}, 1, x_{j+1}, \ldots, x_d) : x_i \in [n_i], i \in [d] - \{j\}, \sum_{i \in [d] - \{j\}} i \cdot x_i \equiv 0 \pmod{2d+1}\}, \]
and let $C_j$ be the set of vertices

$$C_j := \{(x_1, \ldots, x_{j-1}, n_j, x_{j+1}, \ldots, x_d) : x_i \in [n_i], i \in [d] - \{j\}, \sum_{i \in [d] - \{j\}} i \cdot x_i \equiv -j(n_j + 1) \pmod{2d + 1}\}.$$ 

Let $T := \bigcup_j (S \cup B_j \cup C_j)$. Then $T$ is dominating in $P_{n_1} \square P_{n_2} \square \cdots \square P_{n_d}$.

Proof. Consider a vertex $x = (x_1, x_2, \ldots, x_d)$ not in $S$. We now prove that $x$ has neighbour in $S$, or $x$ is in some $B_j \cup C_j$. Now $x_i \in [n_i]$ for each $i \in [d]$, and for some $r \in [2d]$,

$$\sum_{i=1}^{d} i \cdot x_i \equiv r \pmod{2d + 1}.$$ 

First suppose that $r \in [d]$. Let $j := r$. Thus

$$j \cdot (x_j - 1) + \sum_{i \in [d] - \{j\}} i \cdot x_i \equiv 0 \pmod{2d + 1}.$$ 

Hence, if $x_j \neq 1$ then $(x_1, \ldots, x_{j-1}, x_j - 1, x_{j+1}, \ldots, x_d)$ is a neighbour of $x$ in $S$, and $x$ is dominated. If $x_j = 1$ then $x$ is in $B_j \subset T$.

Now assume that $r \in [d + 1, 2d]$. Let $j := 2d + 1 - r \in [d]$. Thus $r \equiv -j \pmod{2d + 1}$, and

$$j \cdot (x_j + 1) + \sum_{i \in [d] - \{j\}} i \cdot x_i \equiv 0 \pmod{2d + 1}.$$ 

Hence, if $x_j \neq n_j$ then $(x_1, \ldots, x_{j-1}, x_j + 1, x_{j+1}, \ldots, x_d)$ is a neighbour of $x$ in $S$, and $x$ is dominated. If $x_j = n_j$ then $x$ is in $C_j \subset T$.

Thus every vertex not in $T$ has a neighbour in $S \subset T$, and $T$ is dominating. \qed

We now determine the size of the dominating set in Lemma 6.1.

Lemma 6.2. For integers $r \geq 2$, $d \geq 1$, $c$, and $n_1, \ldots, n_d \geq 1$, define

$$Q(n_1, \ldots, n_d; c; r) := \{(x_1, x_2, \ldots, x_d) : x_i \in [n_i], i \in [d], \sum_{i \in [d]} i \cdot x_i \equiv c \pmod{r}\}.$$ 

If each $n_i \equiv 0 \pmod{r}$ then for every integer $c$,

$$|Q(n_1, \ldots, n_d; c; r)| = \frac{1}{r} \prod_{i \in [d]} n_i.$$ 

Proof. We proceed by induction on $d$. First suppose that $d = 1$. Without loss of generality, $c \in [r]$. Then

$$Q(n_1; c; r) = \{x \in [n_1] : x \equiv c \pmod{r}\} = \{r \cdot y + c : y \in [0, \frac{n_1}{r} - 1]\}.$$
Thus \(|Q(n_1; c; r)| = \frac{n_1}{r}\), as desired. Now assume that \(d \geq 2\). Thus
\[
|Q(n_1, \ldots, n_d; c; r)| = |\{(x_1, x_2, \ldots, x_d) : x_i \in [n_i], i \in [d], \sum_{i \in [d]} i \cdot x_i \equiv c \pmod{r}\}|
\]
\[= \sum_{x_d \in [n_d]} |\{(x_1, x_2, \ldots, x_d) : x_i \in [n_i], i \in [d - 1], \sum_{i \in [d - 1]} i \cdot x_i \equiv (c - d \cdot x_d) \pmod{r}\}|
\]
\[= \sum_{x_d \in [n_d]} |Q(n_1, \ldots, n_{d-1}; c - d \cdot x_d; r)| .
\]

By induction,
\[|Q(n_1, \ldots, n_d; c; r)| = \sum_{x_d \in [n_d]} \frac{1}{r} \prod_{i \in [d - 1]} n_i = \frac{1}{r} \prod_{i \in [d]} n_i ,
\]
as desired. \(\square\)

**Lemma 6.3.** Let \(G := P_{n_1} \square P_{n_2} \square \cdots \square P_{n_d}\) for some integers \(d \geq 1\) and \(n_1, n_2, \ldots, n_d \geq 1\), where each \(n_i \equiv 0 \pmod{2^d + 1}\). Then
\[
\gamma(G) \leq \frac{\nu(G)}{2d + 1} \left(1 + \sum_{j \in [d]} \frac{2}{n_j}\right).
\]

**Proof.** Using the notation in Lemma 6.1 by Lemma 6.2 applied three times,
\[
|S| = \frac{1}{2d + 1} \prod_{i \in [d]} n_i = \frac{\nu(G)}{2d + 1} ,
\]
and for each \(j \in [d],\)
\[
|B_j|, |C_j| = \frac{1}{2d + 1} \prod_{i \in [d] - \{j\}} n_i = \frac{\nu(G)}{(2d + 1)n_j} .
\]
Thus
\[
|T| \leq |S| + \sum_{j \in [d]} |B_j| + |C_j| = \frac{\nu(G)}{2d + 1} \left(1 + \sum_{j \in [d]} \frac{2}{n_j}\right) ,
\]
as desired. \(\square\)

From the dominating set given in Lemma 6.3 we construct a connected dominating set as follows.

**Lemma 6.4.** Let \(G := P_{n_1} \square P_{n_2} \square \cdots \square P_{n_d}\) for some integers \(d \geq 1\) and \(n_1 \geq n_2 \geq \cdots \geq n_d \geq 1\), where each \(n_i \equiv 0 \pmod{2d + 1}\). Then
\[
\gamma_c(G) < \frac{\nu(G)}{2d - 1} \left(1 + \frac{2d - 2}{n_d} + \sum_{j \in [d - 1]} \frac{2}{n_j}\right) .
\]
Proof. Let $G' := P_{n_1} \Box P_{n_2} \Box \cdots \Box P_{n_d}$. By Lemma 6.3 applied to $G'$,
\[
\gamma(G') \leq \frac{\nu(G')}{{2d} - 1} \left(1 + \sum_{j \in [d-1]} \frac{2}{n_j}\right).
\]
By Lemma 5.8 with $H = G'$,
\[
\gamma_c(G) \leq (n_d - 1) \cdot \gamma(G') + \nu(G').
\]
Thus,
\[
\gamma_c(G) \leq (n_d - 1) \frac{\nu(G')}{{2d} - 1} \left(1 + \sum_{j \in [d-1]} \frac{2}{n_j}\right) + \nu(G')
\]
\[
= \frac{\nu(G)}{{2d} - 1} \left(1 + \sum_{j \in [d-1]} \frac{2}{n_j}\right) - \frac{\nu(G')}{2d - 1} \left(1 + \sum_{j \in [d-1]} \frac{2}{n_j}\right) + \nu(G')
\]
\[
< \frac{\nu(G)}{{2d} - 1} \left(1 + \sum_{j \in [d-1]} \frac{2}{n_j}\right) - \frac{\nu(G')}{{2d} - 1} + \nu(G')
\]
\[
= \frac{\nu(G)}{{2d} - 1} \left(1 + \sum_{j \in [d-1]} \frac{2}{n_j}\right) + \frac{\nu(G')}{{2d} - 1} \left(2d - 2\right)
\]
\[
= \frac{\nu(G)}{{2d} - 1} \left(1 + \frac{2d - 2}{n_d} + \sum_{j \in [d-1]} \frac{2}{n_j}\right),
\]
as desired. \qed

Note that Gravier [18] proved an analogous result to Lemma 6.4 for the total domination number of multi-dimensional grids. Lemma 5.4 leads to the following bounds on the Hadwiger number of even-dimensional grids. These lower and upper bounds are within a multiplicative factor of approximately $2\sqrt{d}$, ignoring lower order terms.

**Theorem 6.5.** Let $G := P_{n_1}^2 \Box P_{n_2}^2 \Box \cdots \Box P_{n_d}^2$ for some integers $d \geq 1$ and $n_1 \geq n_2 \geq \cdots \geq n_d \geq 1$, where each $n_i \equiv 0 \pmod{2d + 1}$. Then
\[
\sqrt{\nu(G)} \left(1 - \frac{1}{{2d} - 1}\right) \left(1 - \frac{1}{d - 1} \sum_{j \in [d-1]} \frac{1}{n_j}\right) + 2 \leq \eta(G) \leq \sqrt{(4d - 2) \cdot \nu(G) + 3}.
\]

**Proof.** The upper bound follows from Lemma 2.2(c) since $\Delta(G) = 4d$. For the lower bound, let $G' := P_{n_1} \Box P_{n_2} \Box \cdots \Box P_{n_d}$. By Lemma 6.4 applied to $G'$,
\[
\gamma_c(G') < \frac{\nu(G')}{2d - 1} \left(1 + \frac{2d - 2}{n_d} + \sum_{j \in [d-1]} \frac{2}{n_j}\right).
\]
By Corollary 5.7 applied to $G'$,
\[
\eta(G) \geq \nu(G') - \gamma_c(G') + 2.
\]
Thus
\[ \eta(G) \geq v(G') - \frac{v(G')}{2d-1} \left( 1 + \frac{2d-2}{n_d} + \sum_{j \in [d-1]} \frac{2}{n_j} \right) + 2 \]
\[ = v(G') \left( 1 - \frac{1}{2d-1} - \frac{2d-2}{(2d-1)n_d} - \frac{1}{2d-1} \sum_{j \in [d-1]} \frac{2}{n_j} \right) + 2 \]
\[ = \sqrt{v(G)} \left( \frac{2d-2}{2d-1} - \frac{2d-2}{(2d-1)n_d} - \frac{2}{2d-1} \sum_{j \in [d-1]} \frac{1}{n_j} \right) + 2 \]
\[ = \sqrt{v(G)} \left( 1 - \frac{1}{2d-1} \right) \left( 1 - \frac{1}{n_d} - \frac{1}{d-1} \sum_{j \in [d-1]} \frac{1}{n_j} \right) + 2 , \]
as desired. \qed

**Corollary 6.6.** For all even integers \( d \geq 4 \) and \( n \geq 1 \) such that \( n \equiv 0 \pmod{2d+1} \),
\[ n^{d/2} \left( 1 - \frac{1}{d-1} \right) \left( 1 - \frac{2}{n} \right) + 2 \leq \eta(P_n^d) \leq \sqrt{2}(d-1)n^{d/2} + 3 . \]

7. Hadwiger Number of Products of Complete Graphs

In this section we consider the Hadwiger number of the product of complete graphs. First consider the case of two complete graphs. Chandran and Raju [5, 38] proved that \( \eta(K_n \boxtimes K_m) = \Theta(n\sqrt{m}) \) for \( n \geq m \). In particular,
\[ \frac{1}{4}(n - \sqrt{m})(\sqrt{m} - 2) \leq \eta(K_n \boxtimes K_m) \leq 2n\sqrt{m} . \]
Since \( P_{\lfloor \sqrt{m} \rfloor} \boxtimes P_{\lfloor \sqrt{m} \rfloor} \boxtimes K_n \) is a subgraph of \( K_n \boxtimes K_m \), Lemma 10.6 below immediately improves this lower bound to
\[ \eta(K_n \boxtimes K_m) \geq \left\lfloor \frac{n}{2} \right\rfloor \lfloor \sqrt{m} \rfloor . \]
Moreover, the proof of Lemma 10.6 is much simpler than the proof in [5, 38].

In the following theorem we sharpen this lower bound, determining \( \eta(K_n \boxtimes K_m) \) to within a factor of \( 2\sqrt{2} \) (ignoring lower order terms).

**Theorem 7.1.** For all integers \( n \geq m \geq 1 \),
\[ n^{\frac{m}{2}} \leq \theta(n + \sqrt{m}) \leq \eta(K_n \boxtimes K_m) < \sqrt{(n + m - 2)(nm + 3} \leq n\sqrt{2m} + 3 . \]

**Proof.** The upper bound follows from Lemma 2.2(c) since \( K_n \boxtimes K_m \) has \( nm \) vertices and maximum degree \( n + m - 2 \).

For the lower bound, let \( p := \left\lfloor \frac{n}{2} \right\rfloor \) and \( k := \left\lfloor \frac{1}{2}(\sqrt{8m+1} - 1) \right\rfloor \). Observe that \( n \geq 2p \) and \( m \geq \frac{1}{4}k(k+1) \). Thus it suffices to prove that the desired minor is in \( K_{2p} \boxtimes K_{k(k+1)/2} \).

Let \( V(K_{2p}) = [2p] \) and \( V(K_{k(k+1)/2}) = \{ (x, y) : 1 \leq x \leq y \leq k \} \). Each vertex is described by a triple \( (j, x, y) \) where \( j \in [2p] \) and \( 1 \leq x \leq y \leq k \). Distinct vertices \( (j, x, y) \) and \( (j', x', y') \) are adjacent if and only if \( x = x' \) and \( y = y' \), or \( j = j' \).

For \( i \in [k] \) and \( j \in [p] \), let \( T_{i,j} \) be the subgraph induced by
\[ \{ (2j-1, x, i) : x \in [i] \} \cup \{ (2j, i, y) : y \in [i, k] \} . \]
Observe that $T_{i,j}$ consists of two cliques connected by the edge $(2j-1, i, i)(2j, i, i)$. Thus $T_{i,j}$ is connected. For $i \leq a$, each pair of distinct subgraphs $T_{i,j}$ and $T_{a,b}$ are disjoint, and are connected by an edge from $(2j, i, a)$ in $T_{i,j}$ to $(2b - 1, i, a)$ in $T_{a,b}$.

Hence the $T_{i,j}$ are branch sets of a $K_{kp}$-minor. Therefore

$$\eta(K_{2p} \square K_{k(k+1)/2}) \geq kp = \left\lfloor \frac{1}{2}(\sqrt{8m + 1} - 1) \right\rfloor \geq n\sqrt{\frac{m}{kn}} - O(n + \sqrt{m}),$$

as desired.

Theorem 7.1 is improved for small values of $m$ as follows.

**Proposition 7.2.** For every integer $n \geq 1$,

$$\eta(K_n \square K_2) = n + 1.$$

**Proof.** Say $V(K_n) = [n]$ and $V(K_2) = \{v, w\}$.

First we prove the lower bound $\eta(K_n \square K_2) \geq n + 1$. For $i \in [n]$, let $X_i$ be the subgraph of $K_n \square K_2$ induced by the vertex $(i, v)$. Let $X_{n+1}$ be the subgraph of $K_n \square K_2$ induced by the vertices $(1, w), \ldots, (n, w)$. Then $X_1, \ldots, X_{n+1}$ are branch sets of a $K_{n+1}$-minor in $K_n \square K_2$. Thus $\eta(K_n \square K_2) \geq n + 1$.

It remains to prove the upper bound $\eta(K_n \square K_2) \leq n + 1$. Let $X_1, \ldots, X_k$ be the branch sets of a complete minor in $K_n \square K_2$, where $k = \eta(K_n \square K_2)$. If every $X_i$ has at least two vertices then $k \leq n$ since $K_n \square K_2$ has $2n$ vertices. Otherwise some $X_i$ has only one vertex, which has degree $n$ in $K_n \square K_2$. Thus $k \leq n + 1$, as desired. □

**Proposition 7.3.** For every integer $n \geq 1$,

$$\eta(K_n \square K_3) = n + 2.$$

**Proof.** A $K_{n+2}$-minor in $K_n \square K_3$ is obtained by contracting the first row and contracting the second row. Thus $\eta(K_n \square K_3) \geq n + 2$.

It remains to prove the upper bound $\eta(K_n \square K_3) \leq n + 2$. We proceed by induction on $n$. The base case $n = 1$ is trivial. Let $X_1, \ldots, X_k$ be the branch sets of a $K_k$-minor, where $k = \eta(K_n \square K_3)$. Without loss of generality, each $X_i$ is an induced subgraph.

Suppose that some column $C$ intersects at most one branch set $X_i$. Deleting $C$ and $X_i$ gives a $K_{k-1}$-minor in $K_{n-1} \square K_3$. By induction, $k - 1 \geq n + 1$. Thus $k \geq n + 2$, as desired. Now assume that every column intersects at least two branch sets.

If some branch set has only one vertex $v$, then $k \leq 1 + \deg(v) = n + 2$, as desired. Now assume that every branch set has at least two vertices.

Suppose that some branch set $X_i$ has vertices in distinct rows. Since $X_i$ is connected, $X_i$ has at least two vertices in some column $C$. Now $C$ intersects at least two branch sets, $X_i$ and $X_j$. Thus $X_j$ intersects $C$ in exactly one vertex $v$. Consider the subgraph $X_j - v$. It has at least one vertex. Every neighbour of $v$ that is in $X_j$ is in the same row as $v$. Since $X_j$ is an induced subgraph, the neighbourhood of $v$ in $X_j$ is a non-empty clique. Thus $v$ is not a cut-vertex in $X_j$, and $X_j$ is a non-empty connected subgraph. Hence deleting $C$ and $X_i$ gives a $K_{k-1}$-minor in $K_{n-1} \square K_3$. By induction, $k - 1 \geq n + 1$. Thus $k \geq n + 2$, as desired. Now assume that each branch set is contained in some row.
If every branch set has at least three vertices, then \( k \leq \frac{1}{3!} |V(K_n \square K_3)| = n \), as desired. Now assume that some branch set \( X_i \) has exactly two vertices \( v \) and \( w \). Now \( v \) and \( w \) are in the same row. There are at most \( \frac{n-k}{2} \) other branch sets in the same row, since every branch set has at least two vertices and is contained in some row. Moreover, \( N(v) \cup N(w) \) contains only four vertices that are not in the same row as \( v \) and \( w \). Thus \( k - 1 \leq \frac{n-k}{2} + 4 \). That is, \( k \leq \frac{n-k+8}{2} \), which is at most \( n+2 \) whenever \( n \geq 4 \). Now assume that \( n \leq 3 \). Since every branch set has at least two vertices, \( k \leq \frac{1}{3!} |V(K_n \square K_3)| = \frac{3p+4}{2} \), which is at most \( n+2 \) for \( n \leq 4 \).

**Proposition 7.4.** For every integer \( n \geq 1 \),

\[
3 \left\lceil \frac{n}{2} \right\rceil \leq \eta(K_n \square K_4) \leq 2n + 3.
\]

**Proof.** The upper bound follows from Lemma 2.2(c) since \( \nu(K_n \square K_4) = 4n \) and \( \Delta(K_n \square K_4) = n + 2 \).

For the lower bound, let \( p := \left\lceil \frac{n}{2} \right\rceil \). Since \( n \geq 2p \), it suffices to prove that the desired minor is in \( K_{2p} \square K_4 \). Each vertex is described by a pair \((i, x)\) where \( i \in [2p] \) and \( x \in \{a, b, c, d\} \). Distinct vertices \((i, x)\) and \((j, y)\) are adjacent if and only if \( i = j \) or \( x = y \).

For \( i \in [p] \), let \( X_i \) be the path \((2i-1, a)(2i-1, b)(2i, b)(2i, c)\), let \( Y_i \) be the edge \((2i-1, c)(2i, d)\), and let \( Z_i \) be the edge \((2i, a)(2i, d)\). Thus each \( X_i, Y_i, \) and \( Z_i \) is connected, and each pair of distinct subgraphs are disjoint. Moreover, the vertex \((2i-1, a)\) in \( X_i \) is adjacent to the vertex \((2j-1, a)\) in \( X_j \). The vertex \((2i, c)\) in \( X_i \) is adjacent to the vertex \((2j-1, c)\) in \( Y_j \). The vertex \((2i-1, a)\) in \( X_i \) is adjacent to the vertex \((2j, a)\) in \( Z_j \). The vertex \((2i-1, c)\) in \( Y_i \) is adjacent to the vertex \((2j-1, c)\) in \( Y_j \). The vertex \((2i-1, d)\) in \( Y_i \) is adjacent to the vertex \((2j, d)\) in \( Z_j \). And the vertex \((2i, a)\) in \( Z_i \) is adjacent to the vertex \((2j, a)\) in \( Z_j \). Hence \( \{X_i, Y_i, Z_i : i \in [p]\} \) are the branch sets of a \( K_{2p} \)-minor. Therefore \( \eta(K_{2p} \square K_4) \geq 3p \).

We conjecture that the upper bound in Proposition 7.4 can be improved to \( \eta(K_n \square K_m) \leq (2 - \epsilon)n \) for some \( \epsilon > 0 \) and for all sufficiently large \( n \).

Now consider the Hadwiger number of the product of \( d \) complete graphs. Here our lower and upper bounds are within a factor of \( 2\sqrt{d} \) (ignoring lower order terms).

**Theorem 7.5.** For all integers \( n_1 \geq n_2 \geq \cdots \geq n_d \geq 2 \),

\[
\left\lceil \frac{n_1}{2} \right\rceil \prod_{i \in [2, d]} \sqrt{n_i} \leq \eta(K_{n_1} \square K_{n_2} \square \cdots \square K_{n_d}) < \sqrt{d} n_1 \prod_{i \in [2, d]} \sqrt{n_i} + 3.
\]

**Proof.** Let \( G := K_{n_1} \square K_{n_2} \square \cdots \square K_{n_d} \).

Since \( \nu(G) = \prod_{i} n_i \) and \( \Delta(G) = \sum_i (n_i - 1) \), Lemma 2.2(c) implies the upper bound,

\[
\eta(G) < \sqrt{\left( \sum_{i \in [d]} (n_i - 1) \right) \prod_{i \in [d]} \sqrt{n_i} + 3} = \sqrt{d} n_1 \prod_{i \in [2, d]} \sqrt{n_i} + 3.
\]

For the lower bound, let \( p := \left\lceil \frac{n_1}{2} \right\rceil \) and \( k_i := \left\lceil \sqrt{n_i} \right\rceil \) for each \( i \in [2, d] \). Let \( m_1 := 2p \) and \( m_i := k_i^2 \) for each \( i \in [2, d] \). Observe that each \( n_i \geq m_i \). Thus it suffices to construct
the desired minor in \(K_{m_1} \square K_{m_2} \square \cdots \square K_{m_d}\). Let \(V(K_{m_1}) = [2p]\), and for each \(i \in [2, d]\), let
\[
V(K_{m_i}) = \{(a_i, b_i) : a_i, b_i \in [k_i]\}.
\]
Each vertex is described by a vector \((r, a_2, b_2, \ldots, a_d, b_d)\) where \(r \in [m_1]\) and \(a_i, b_i \in [k_i]\).
Distinct vertices \((r, a_2, b_2, \ldots, a_d, b_d)\) and \((s, x_2, y_2, \ldots, x_d, y_d)\) are adjacent if and only if:
(1) \(a_i = x_i\) and \(b_i = y_i\) for each \(i \in [2, d]\), or
(2) \(r = s\), and for some \(i \in [2, d]\), for every \(j \neq i\), we have \(a_j = x_j\) and \(b_j = y_j\).
In case (1) the edge is in dimension \(1\) and in case (2) the edge is in dimension \(i\).

For all \(r \in [p]\), \(i \in [2, d]\), and \(j_i \in [k_i]\), let \(A\langle r, j_2, \ldots, j_d \rangle\) be the subgraph induced by
\[
\{(2r, a_2, j_2, j_3, \ldots, a_d, j_d) : a_i \in [k_i], i \in [2, d]\},
\]
let \(B\langle r, j_2, \ldots, j_d \rangle\) be the subgraph induced by
\[
\{(2r - 1, j_2, b_2, j_3, \ldots, j_d, b_d) : b_i \in [k_i], i \in [2, d]\},
\]
and let \(X\langle r, j_2, \ldots, j_d \rangle\) be the subgraph induced by the vertex set of \(A\langle r, j_2, \ldots, j_d \rangle \cup B\langle r, j_2, \ldots, j_d \rangle\).

Observe that any two vertices in \(A\langle r, j_2, \ldots, j_d \rangle\) are connected by a path of at most \(d - 1\) edges (in dimensions \(2, \ldots, d\)). Thus \(A\langle r, j_2, \ldots, j_d \rangle\) is connected. Similarly, \(B\langle r, j_2, \ldots, j_d \rangle\) is connected. Moreover, the dimension-1 edge
\[
(2r, j_2, j_2, j_3, j_3, \ldots, j_d, j_d)(2r - 1, j_2, j_2, j_3, j_3, \ldots, j_d, j_d)
\]
connects \(A\langle r, j_2, \ldots, j_d \rangle\) and \(B\langle r, j_2, \ldots, j_d \rangle\). Hence \(X\langle r, j_2, \ldots, j_d \rangle\) is connected.

Consider a pair of distinct subgraphs \(X\langle r, j_2, \ldots, j_d \rangle\) and \(X\langle s, \ell_2, \ldots, \ell_d \rangle\). By construction they are disjoint. Moreover, the dimension-1 edge
\[
(2r, \ell_2, j_2, \ell_3, j_3, \ldots, \ell_d, j_d)(2s - 1, \ell_2, j_2, \ell_3, j_3, \ldots, \ell_d, j_d)
\]
connects \(A\langle r, j_2, \ldots, j_d \rangle\) and \(B\langle s, \ell_2, \ldots, \ell_d \rangle\). Hence the \(X\langle r, j_2, \ldots, j_d \rangle\) are branch sets of a clique minor of order \(p \prod_{i=2}^d k_i\). Therefore
\[
\eta(G) \geq p \prod_{i=2}^d k_i = \left\lfloor \frac{n_1}{2} \right\rfloor \prod_{i=2}^d \left\lfloor \sqrt{n_i} \right\rfloor,
\]
as desired. \(\square\)

The \(d\)-dimensional Hamming graph is the product
\[
H_n^d := K_n \square K_n \square \cdots \square K_n.
\]

Chandran and Sivadasan \([7]_{\text{a}}\) proved the following bounds on the Hadwiger number of \(H_n^d\):
\[
n^{(d-1)/2} \leq \eta(H_n^d) \leq 1 + \sqrt{d} n^{(d+1)/2}.
\]
Theorem 7.5 improves this lower bound by a \(\Theta(n)\) factor; thus determining \(\eta(H_n^d)\) to within a \(2\sqrt{d}\) factor (ignoring lower order terms):
\[
\frac{1}{2} n^{(d+1)/2} - \mathcal{O}(n^{d/2}) \leq \eta(H_n^d) < 1 + \sqrt{d} n^{(d+1)/2}.
\]
8. Hypercubes and Lexicographic Products

The \(d\)-dimensional hypercube is the graph

\[ Q_d := K_2 \square K_2 \square \ldots \square K_2. \]

Hypercubes are both grid graphs and Hamming graphs. The Hadwiger number of \(Q_d\) was first studied by Chandran and Sivadasan \([3,2]\). The best bounds on \(\eta(Q_d)\) are due to Kotlov \([32]\), who proved that

\[
\eta(Q_d) \geq \begin{cases} 
2^{(d+1)/2}, & d \text{ odd} \\
3 \cdot 2^{(d-2)/2}, & d \text{ even}
\end{cases}
\]

and

\[
\eta(Q_d) \leq \frac{5}{2} + \sqrt{2^d (d - 3) + \frac{25}{4}}.
\]

Kotlov \([32]\) actually proved the following more general result which readily implies (3) by induction:

**Proposition 8.1 ([32]).** For every bipartite graph \(G\), the strong product \(G \square K_2\) is a minor of \(G \square K_2 \square K_2\).

Proposition 8.1 is generalised by the following result with \(H = K_2\) (since \(G \cdot K_2 \cong G \square K_2 \square K_2\)).

**Proposition 8.2.** For every bipartite graph \(G\) and every connected graph \(H\), the lexicographic product \(G \cdot H\) is a minor of \(G \square H \square H\).

**Proof.** Properly colour the vertices of \(G\) black and white. For all vertices \((v, p)\) of \(G \cdot H\), let \(X\langle v, p \rangle\) be the subgraph of \(G \square H \square H\) induced by the set

\[
\begin{cases} 
\{(v, p, q) : q \in V(H)\}, & \text{if } v \text{ is black} \\
\{(v, q, p) : q \in V(H)\}, & \text{if } v \text{ is white}.
\end{cases}
\]

We claim that the \(X\langle v, p \rangle\) form the branch sets of a \(G \cdot H\)-minor in \(G \square H \square H\). First observe that each \(X\langle v, p \rangle\) is isomorphic to \(H\), and is thus connected.

Consider distinct vertices \((v, p)\) and \((v', p')\) of \(G \cdot H\), where \(v\) is black.

Suppose on the contrary that some vertex \((w, a, b)\) of \(G \square H \square H\) is in both \(X\langle v, p \rangle\) and \(X\langle v', p' \rangle\). Since \((w, a, b)\) is in \(X\langle v, p \rangle\), we have \(a = p\). By construction, \(w = v = v'\). Thus \(v'\) is also black, and since \((w, a, b)\) is in \(X\langle v', p' \rangle\), we have \(a = p'\). Hence \(p = p'\), which contradicts that \((v, p)\) and \((v', p')\) are distinct. Hence \(X\langle v, p \rangle\) and \(X\langle v', p' \rangle\) are disjoint.

Suppose that \((v, p)\) and \((v', p')\) are adjacent in \(G \cdot H\). It remains to prove that \(X\langle v, p \rangle\) and \(X\langle v', p' \rangle\) are adjacent in \(G \square H \square H\). By definition, \(vv' \in E(G)\), or \(v = v'\) and \(pp' \in E(H)\). If \(vv' \in E(G)\), then without loss of generality, \(v\) is black and \(v'\) is white, implying that \((v, p, p')\), which is in \(X\langle v, p \rangle\), is adjacent to \((v', p, p')\), which is in \(X\langle v', p' \rangle\). If \(v = v'\) and \(pp' \in E(H)\), then for every \(q \in V(H)\), the vertex \((v, p, q)\), which is in \(X\langle v, p \rangle\), is adjacent to \((v, p', q)\), which is in \(X\langle v', p' \rangle\). \(\square\)
Proposition 8.2 motivates studying \( \eta(G \cdot H) \). We now show that when \( G \) is a complete graph, \( \eta(G \cdot H) \) can be determined precisely.

**Proposition 8.3.** For every graph \( H \),

\[
\eta(K_n \cdot H) = \left\lfloor \frac{n}{2} (\nu(H) + \omega(H)) \right\rfloor.
\]

*Proof.* Let \( C \) be a maximum clique in \( H \). Consider the set of vertices

\[ X := \{(u, y) : u \in V(K_n), y \in C\} \]

in \( K_n \cdot H \). Thus \( |X| = n \cdot \omega(H) \) and \( X \) is a clique in \( K_n \cdot H \). In fact, \( X \) is a maximum clique, since every set of \( n \cdot \omega(H) + 1 \) vertices in \( K_n \cdot H \) contains \( \omega(H) + 1 \) vertices in a single copy of \( H \). Thus \( \omega(K_n \cdot H) = n \cdot \omega(H) \). Hence the upper bound on \( \eta(K_n \cdot H) \) follows from Lemma 2.3.

It remains to prove the lower bound on \( \eta(K_n \cdot H) \). Delete the edges of \( H \) that are not in \( C \). This operation is allowed since it does not increase \( \eta(K_n \cdot H) \). So \( H \) now consists of \( C \) and some isolated vertices. Observe that \( (K_n \cdot H) \setminus X \) is isomorphic to the balanced complete \( n \)-partite graph with \( \nu(H) - \omega(H) \) vertices in each colour class. Every balanced complete multipartite graph with \( r \) vertices has a matching of \( \left\lfloor \frac{r}{2} \right\rfloor \) edges [47]. Thus \( (K_n \cdot H) \setminus X \) has a matching \( M \) of \( \left\lfloor \frac{n}{2} (\nu(H) - \omega(H)) \right\rfloor \) edges. No edge in \( M \) is incident to a vertex in \( X \). For every edge \( (v, x)(v', x') \) in \( M \) and vertex \( (u, y) \) of \( K_n \cdot H \), since \( v \neq v' \), without loss of generality, \( v \neq u \). Thus \( vu \in E(K_n) \) and \( (v, x) \) is adjacent to \( (u, y) \) in \( K_n \cdot H \). Hence contracting each edge in \( M \) gives a \( K_{|X| + |M|} \)-minor in \( K_n \cdot H \). Therefore

\[
\eta(K_n \cdot H) \geq |X| + |M| = n \cdot \omega(H) + \left\lfloor \frac{n}{2} (\nu(H) - \omega(H)) \right\rfloor = \left\lfloor \frac{n}{2} (\nu(H) + \omega(H)) \right\rfloor.
\]

We now show that Propositions 8.2 and 8.3 are closely related to some previous results in the paper.

Proposition 8.2 with \( G = K_2 \) implies that \( K_2 \cdot H \) is a minor of \( H \square H \square K_2 \). Proposition 8.3 implies that \( \eta(K_2 \cdot H) = \nu(H) + \omega(H) \). Thus \( \eta(H \square H \square K_2) \geq \nu(H) + \omega(H) \), which is only slightly weaker than Proposition 3.1.

Proposition 8.2 with \( G = K_{n,n} \) implies that \( K_{n,n} \cdot H \) is a minor of \( K_{n,n} \square H \square H \) for every connected graph \( H \). Since \( K_{n+1} \) is a minor of \( K_{n,n} \), we have \( K_{n+1} \cdot H \) is a minor of \( K_{n,n} \square H \square H \). Proposition 8.3 implies that

\[
\eta(K_{n,n} \square H \square H) \geq \left\lfloor \frac{n+1}{2} (\nu(H) + \omega(H)) \right\rfloor.
\]

Since \( K_{n,n} \subset K_{2n} \),

\[
\eta(K_{2n} \square H \square H) \geq \left\lfloor \frac{n+1}{2} (\nu(H) + \omega(H)) \right\rfloor.
\]
With $H = K_d^m$ we have
\[
\eta(K_{2n} \Box K_{2d}^m) \geq \left\lfloor \frac{n + 1}{2} (m^d + m) \right\rfloor,
\]
which is equivalent to Theorem 7.5 with $n_1 = n$ and $n_2 = \cdots = n_{2d+1} = m$ (ignoring lower order terms). In fact for small values of $m$, this bound is stronger than Theorem 7.5. For example, with $n = 1$ and $m = 3$ we have
\[
\eta(K_2 \Box K_{2d}^3) \geq 3d + 3,
\]
whereas Theorem 7.5 gives no non-trivial bound on $\eta(K_2 \Box K_{2d}^3)$.

9. ROUGH STRUCTURAL CHARACTERISATION THEOREM FOR TREES

In this section we characterise when the product of two trees has a large clique minor. Section 5 gives such an example: Corollaries 5.3 and 5.6 imply that if one tree has many leaves and the other has many vertices then their product has a large clique minor. Now we give a different example. As illustrated in Figure 4(a), let $B_n$ be the tree obtained from the path $P_{2n+1}$ by adding one leaf adjacent to the vertex in the middle of $P_{2n+1}$.

Now $B_n$ only has three leaves, but Seese and Wessel [45] implicitly proved that the product of $B_n$ and a long path (which only has two leaves) has a large clique minor. In Theorem 9.1 below we prove an explicit bound of $\eta(P_m \Box B_n) \geq \min\{n, \sqrt{m}\}$, which is illustrated in Figure 5. In fact, this bound holds for a more general class of trees than $B_n$, which we now introduce.

As illustrated in Figure 4(b), a balance of order $n$ is a tree $T$ that has an edge $rs$, and disjoint sets $A, B \subseteq V(T) - \{r, s\}$, each with at least $n$ vertices, such that $A \cup \{r\}$ and $B \cup \{r\}$ induce connected subtrees in $T$. We say $A$ and $B$ are the branches, $r$ is the root, and $s$ is the support of $T$. For example, the star $S_t$ is a balance of order $\left\lfloor \frac{t-1}{2} \right\rfloor$, and $B_n$ is a balance of order $n$. Theorem 9.1 below implies that $\eta(P_m \Box T) \geq \min\{\sqrt{m}, n\}$ for balanced graphs $T$.
Figure 5. $K_n$ is a minor of $P_{n2} \Box B_n$.

every balance $T$ of order $n$. The critical property of a long path is that it has many large disjoint subpaths.

**Theorem 9.1.** Let $G$ be a tree that has $n$ disjoint subtrees each of order at least $n$. Then for every balance $T$ of order $n$,

$$\eta(G \Box T) \geq n .$$

**Proof.** Let $A$ and $B$ be the branches, let $r$ be the root, and let $s$ be the support of $T$. Contract edges in $T$ until $A$ and $B$ each have exactly $n$ vertices. Orient the edges of $T$ away from $r$. Label the vertices of $A$ by $\{a_1, a_2, \ldots, a_n\}$ such that if $\overrightarrow{a_i a_j} \in E(A)$ then $i < j$. Label the vertices of $B$ by $\{b_1, b_2, \ldots, b_n\}$ such that if $\overrightarrow{b_i b_j} \in E(B)$ then $j < i$. For each $i \in [n]$, let $T_i$ be the path between $a_i$ and $b_i$ in $T$ (which thus includes $r$).

Contract edges in $G$ until it is the union of $n$ disjoint subtrees $G_1, \ldots, G_n$, each with exactly $n$ vertices. For every pair of vertices $v, w \in V(G)$, let $G(v, w)$ be the path between $v$ and $w$ in $G$. Orient the edges of $G$ away from an arbitrary vertex in $G_1$. Let $G^*$ be the oriented tree obtained from $G$ by contracting each $G_i$ into a single vertex $z_i$. Note that for each $i \in [2, n]$, each vertex $z_i$ has exactly one incoming arc in $G^*$. Fix a proper 2-colouring of $G^*$ with colours black and white, where $z_1$ is coloured white. Label the vertices in each subtree $G_i$ by $\{v_{i,1}, \ldots, v_{i,n}\}$, such that for each arc $\overrightarrow{(v_{i,j}, v_{i,\ell})}$ in $G_i$, we have $\ell < j$ if $z_i$ is black, and $j < \ell$ if $z_i$ is white.

For each $i \in [n]$, let $H_i$ be the subgraph of $G \Box T$ induced by

$$\{(v_{i,j}, s) : j \in [n]\} .$$

For all $i, j \in [n]$, let $T_{i,j}$ be the subgraph of $G \Box T$ induced by

$$\{(v_{j,i}, y) : y \in V(T_i)\} .$$
For all \(i, j \in [n]\), let \(c_{i,j}\) be the vertex \(a_i\) if \(z_j\) is white, and \(b_i\) if \(z_j\) is black. For all \(i \in [n]\) and \(j \in [2, n]\), let \(U_{i,j}\) be the subgraph of \(G \square T\) induced by
\[
\{(x, c_{i,j}) : x \in G(v_{k,i}, v_{j,i})\},
\]
where \((z_k, z_j)\) is the incoming arc at \(z_j\) in \(G^*\).

For each \(i \in [n]\), let \(X_i\) be the subgraph of \(G \square T\) induced by
\[
\bigcup_{j \in [n]} (T_{i,j} \cup U_{i,j} \cup H_i).
\]
We now prove that \(X_1, \ldots, X_n\) are the branch sets of a \(K_n\)-minor.

First we prove that each \(X_i\) is connected. Observe that each \(H_i\) is isomorphic to \(G_i\), and is thus connected. Each \(T_{i,j}\) is isomorphic to the path \(T_i\), and is thus connected. Moreover, the endpoints of \(T_{i,j}\) are \((v_{j,i}, a_i)\) and \((v_{j,i}, b_i)\). Each \(U_{i,j}\) is isomorphic to the path \(G(v_{j,i}, v_{j,i})\), and is thus connected. Moreover, the endpoints of \(U_{i,j}\) are \((v_{j,i}, c_{i,j})\) and \((v_{j,i}, c_{i,j})\). Thus if \(z_j\) is white (and thus \(z_{j'}\) is black), then the endpoints of \(U_{i,j}\) are \((v_{j,i}, a_i)\) and \((v_{j,i}, a_i)\). And if \(z_j\) is black (and thus \(z_{j'}\) is white), then the endpoints of \(U_{i,j}\) are \((v_{j,i}, b_i)\) and \((v_{j,i}, b_i)\). Hence the induced subgraph
\[
T_{i,1} \cup U_{i,1} \cup T_{i,2} \cup U_{i,2} \cup T_{i,3} \cup U_{i,3} \cup \cdots \cup T_{i,n} \cup U_{i,n}
\]
is a path (illustrated in Figure 5 by alternating vertical and horizontal segments). Furthermore, the vertex \((v_{i,i}, s)\) in \(H_i\) is adjacent to the vertex \((v_{i,i}, r)\) in \(T_{i,i}\). Thus \(X_i\) is connected.

Now we prove that the subgraphs \(X_i\) and \(X_{i'}\) are disjoint for all distinct \(i, i' \in [n]\). First observe that \(H_i\) and \(H_{i'}\) are disjoint since the first coordinate of every vertex in \(H_i\) is some \(v_{i,i}\). Similarly, for all \(j, j' \in [n]\), the subgraphs \(T_{i,j}\) and \(T_{i',j'}\) are disjoint since the first coordinate of every vertex in \(T_{i,j}\) is \(v_{j,i}\). For all \(j, j' \in [n]\), \(U_{i,j}\) and \(U_{i',j'}\) are disjoint since the second coordinate of every vertex in \(U_{i,j}\) is \(a_i\) or \(b_i\). For all \(j \in [n]\), \(H_i\) is disjoint from \(T_{i,j} \cup U_{i,j}\) since the second coordinate of \(H_i\) is \(s\). It remains to prove that \(T_{i,j}\) and \(U_{i',j'}\) are disjoint. Suppose on the contrary that for some \(j, j' \in [n]\), some vertex \((x, y)\) is in \(T_{i,j} \cap U_{i',j'}\). Without loss of generality, \(z_{j'}\) is black. Say \((z_k, z_{j'})\) is the incoming arc at \(z_{j'}\) in \(G^*\). So \(z_k\) is white. Since \((x, y) \in T_{i,j}\), we have \(x = v_{j,i}\) and \(y \in V(T_i)\). Since \((x, y) \in U_{i',j'}, x\) is in the path \(G(v_{k,i'}, v_{j',i'})\). Since \(z_{j'}\) is black, \(y = b_{i'}\). Now \(v_{j,i}\) (which equals \(x\) in the path \(G(v_{k,i'}, v_{j',i'})\)). Thus by the labelling of vertices in \(G_k\) and \(G_{j'}\), we have \(i' < i \leq n\). By comparing the second coordinates, observe that \(b_{i'} \in V(T_i)\). Thus \(i' > i\) by the labelling of the vertices in \(B\). This contradiction proves that \(T_{i,j}\) and \(U_{i',j'}\) are disjoint for all \(j, j' \in [n]\). Hence \(X_i\) and \(X_{i'}\) are disjoint.

Finally, observe that for distinct \(i, i' \in [n]\), the vertex \((v_{i,i'}, s)\) in \(H_i \subset X_i\) is adjacent to the vertex \((v_{i,i'}, r)\) in \(T_{i,i} \subset X_{i'}\). Therefore the \(X_i\) are branch sets of a \(K_n\)-minor.

We conjecture that the construction in Theorem 9.1 is within a constant factor of optimal; that is, \(P_m \square B_n = \Theta(\min\{\sqrt{m}, n\})\).

Theorem 9.1 motivates studying large disjoint subtrees in a given tree. Observe that a star does not have two disjoint subtrees, both with at least two vertices. Thus a star
cannot be used as the tree $G$ in Theorem 9.1 with $n \geq 2$. On the other hand, a path on $n^2$ vertices has $n$ disjoint subpaths, each with $n$ vertices. Of the trees with the same number of vertices, the star has the most leaves and the path has the least. We now prove that the every tree with few leaves has many large disjoint subtrees.\footnote{As an aside we now describe a polynomial-time algorithm that for a given tree $T$, finds the maximum number of disjoint subtrees in $T$, each with at least $n$ vertices. It is convenient to consider a generalisation of this problem, where each vertex $v$ is assigned a positive weight $w(v)$, and each subtree is required to have total weight at least $n$. Let $v$ be a leaf of $T$. Let $T' := T - v$. Define a new weight function $w'(z) := w(z)$ for every vertex $z$ of $T - v$. First suppose that $w(v) \geq n$. Then $T$ has $k$ disjoint subtrees each of total $w$-weight at least $n$ if and only if $T'$ has $k - 1$ disjoint subtrees each of total $w'$-weight at least $n$. (In which case $T[\{v\}]$ becomes one of the $k$ subtrees.) Now assume that $w(v) < n$. Let $x$ be the neighbour of $v$ in $T$. Redefine $w'(x) := w(x) + w(v)$. Then $T$ has $k$ disjoint subtrees each of total $w$-weight at least $n$ if and only if $T'$ has $k$ disjoint subtrees each of total $w'$-weight at least $n$. (A subtree $X$ of $T'$ containing $x$ is replaced by the subtree $T'[V(X) \cup \{v\}]$.)}

**Theorem 9.2.** Let $T$ be a tree with at least one edge, and let $n$ be a positive integer, such that

$$v(T) \geq n^2 + (\text{star}(T) - 2)(n - 1) + 1.$$ 

Then $T$ has $n$ disjoint subtrees, each with at least $n$ vertices.

The proof of Theorem 9.2 is based on the following lemma.

**Lemma 9.3.** Fix a positive integer $n$. Let $T$ be a tree with at least one edge, where each vertex of $T$ is assigned a weight $w(v) \in \mathbb{Z}^+$, such that every leaf has weight at most $n$ and every other vertex has weight 1. Let $W(T) := \sum_{v \in V(T)} w(v)$ be the total weight. Then there is a vertex-partition of $T$ into at least

$$f(T) := \left\lfloor \frac{W(T) - (\text{star}(T) - 2)(n - 1)}{n} \right\rfloor$$

disjoint subtrees, each with total weight at least $n$.

**Proof.** We proceed by induction on $f(T)$. If $f(T) \leq 0$ then there is nothing to prove. First suppose that $f(T) = 1$. Then $W(T) \geq (\text{star}(T) - 2)(n - 1) + n \geq n$ since $\text{star}(T) \geq 2$. Thus $T$ itself has total weight at least $n$, and we are done. Now assume that $f(T) \geq 2$.

Suppose that $T = K_2$ with vertices $x$ and $y$, both of which are leaves. Thus $\text{star}(T) = 2$ and $f(T) = \lfloor \frac{w(x) + w(y)}{n} \rfloor$. Now $f(T) \geq 2$ and $w(x), w(y) \leq n$. Thus $f(T) = 2$ and $w(x) = w(y) = n$. Hence $T[\{x\}]$ and $T[\{y\}]$ is a vertex-partition of $T$ into two disjoint subtrees, each with weight at least $n$, and we are done. Now assume that $v(T) \geq 3$.

Suppose that $w(v) = n$ for some leaf $v$. Let $T' := T - v$. By induction, there is a vertex-partition of $T'$ into $f(T')$ disjoint subtrees, each with total weight at least $n$. These subtrees plus $T[\{v\}]$ are a vertex-partition of $T$ into $1 + f(T')$ disjoint subtrees, each with total weight at least $n$. Now $W(T') = W(T) - n$, and $s(T') \leq \text{star}(T)$ since $v$ is not a leaf in $T'$, and the neighbour of $v$ is the only potential leaf in $T'$ that is not a leaf.
in $T$. Thus
\[ 1 + f(T') = \left\lfloor \frac{n + W(T') - (s(T') - 2)(n - 1)}{n} \right\rfloor \geq \left\lfloor \frac{W(T) - (\text{star}(T) - 2)(n - 1)}{n} \right\rfloor = f(T), \]
and we are done. Now assume that $w(v) \leq n - 1$ for every leaf $v$.

Let $T'$ be the tree obtained from $T$ by deleting each leaf. Since $T \neq K_2$, $T'$ has at least one vertex. Suppose that $T'$ has exactly one vertex. That is, $T$ is a star. Thus $W(T) \leq 1 + \text{star}(T) \cdot (n - 1)$ since every leaf has weight at most $n - 1$. Thus
\[ f(T) = \left\lfloor \frac{W(T) - (\text{star}(T) - 2)(n - 1)}{n} \right\rfloor \leq \left\lfloor \frac{2n - 1}{n} \right\rfloor = 1, \]
which is a contradiction.

Now assume that $T'$ has at least two vertices. In particular, $T'$ has a leaf $v$. Note that the neighbour of $v$ in $T'$ is not a leaf in $T$, as otherwise $T'$ would only have one vertex. Now $v$ is not a leaf in $T$ (since it is in $T'$). Thus $v$ is adjacent to at least one leaf in $T$. Let $x_1, x_2, \ldots, x_d$ be the leaves of $T$ that are adjacent to $v$.

First suppose that $\sum_i w(x_i) \leq n - 1$. Let $T'' := T - \{x_1, \ldots, x_d\}$. Then $v$ is a leaf in $T''$. Redefine $w(v) := 1 + \sum_i w(x_i)$. By induction, there is a vertex-partition of $T''$ into $f(T'')$ disjoint subtrees, each of total weight at least $n$. Observe that $W(T) = W(T'')$ and $s(T'') = \text{star}(T) - d + 1 \leq \text{star}(T)$. Thus
\[ f(T'') = \left\lfloor \frac{W(T'') - (s(T'') - 2)(n - 1)}{n} \right\rfloor \geq \left\lfloor \frac{W(T) - (\text{star}(T) - 2)(n - 1)}{n} \right\rfloor = f(T) . \]
Adding $x_1, \ldots, x_d$ to the subtree of $T''$ containing $v$ gives a vertex-partition of $T$ into at least $f(T)$ disjoint subtrees, each with total weight at least $n$.

Now assume that $\sum_i w(x_i) \geq n$. Let $T''' := T - \{v, x_1, \ldots, x_d\}$. Now $s(T''') \leq \text{star}(T) - d + 1$, since $w_1, \ldots, w_d$ are leaves in $T$ that are not in $T'''$, and the neighbour of $v$ in $T'''$ is the only vertex in $T'''$ that possibly is a leaf in $T'''$ but not in $T$. Observe that $W(T''') \geq W(T) - 1 - d(n - 1)$ since $v$ has weight 1 and each $w_i$ has weight at most $n - 1$. By induction, there is partition of $V(T''')$ into at least $f(T''')$ disjoint subtrees, each with total weight at least $n$. These subtrees plus $T[\{v, x_1, \ldots, x_d\}]$ are a vertex-partition of $T$ into at least $1 + f(T''')$ disjoint subtrees, each with total weight at least $n$. Now
\[ 1 + f(T''') = 1 + \left\lfloor \frac{W(T''') - (s(T''') - 2)(n - 1)}{n} \right\rfloor \geq \left\lfloor \frac{n + W(T) - 1 - d(n - 1) - (\text{star}(T) - d + 1 - 2)(n - 1)}{n} \right\rfloor = \left\lfloor \frac{W(T) - (\text{star}(T) - 2)(n - 1)}{n} \right\rfloor = f(T) . \]
This completes the proof. \qed
Lemma 9.4. For every positive integer $n$, every tree $T$ with at least one edge has

$$\left\lfloor \frac{v(T) - (\text{star}(T) - 2)(n-1)}{n} \right\rfloor$$

disjoint subtrees, each with at least $n$ vertices. Moreover, for all integers $s, n \geq 2$ and $N \geq sn$, there is a tree $T$ with $v(T) = N$ and $\text{star}(T) = s$, such that $T$ has at most

$$\left\lfloor \frac{v(T) - (\text{star}(T) - 2)(n-1)}{n} \right\rfloor$$

disjoint subtrees, each with at least $n$ vertices.

Proof. Lemma 9.3 with each leaf assigned a weight of 1 implies the first claim. It remains to construct the tree $T$. Fix a path $P$ with $N - s(n-1)$ vertices. Let $v$ and $w$ be the endpoints of $P$. As illustrated in Figure 6, let $T$ be the tree obtained from $P$ by attaching $\left\lceil \frac{s}{2} \right\rceil$ pendant paths to $v$, each with $n-1$ vertices, and by attaching $\left\lfloor \frac{s}{2} \right\rfloor$ pendant paths to $w$, each with $n-1$ vertices. Thus $T$ has $N$ vertices and $s$ leaves.

![Figure 6. The tree $T$ in Lemma 9.4](image)

Let $A$ be the subtree of $T$ induced by the union of $v$ and the pendant paths attached at $v$. Let $B$ be the subtree of $T$ induced by the union of $w$ and the pendant paths attached at $w$. Let $X_1, \ldots, X_t$ be a set of disjoint subtrees in $T$, each with at least $n$ vertices. If some $X_i$ intersects $A$ then $v$ is in $X_i$. At most one subtree $X_i$ contains $v$. Thus at most one subtree $X_i$ intersects $A$. Similarly, at most one subtree $X_j$ intersects $B$. The remaining $t-2$ subtrees are contained in $P - \{v, w\}$. Thus

$$t - 2 \leq \frac{N - s(n-1) - 2}{n}.$$ 

It follows that $tn \leq N - (s - 2)(n - 1)$, as desired. \qed

Proof of Theorem 9.2. The assumption in Theorem 9.2 implies that

$$\frac{v(T) - (\text{star}(T) - 2)(n-1)}{n} \geq n + \frac{1}{n}.$$
Hence
\[ \left\lfloor \frac{v(T) - (\text{star}(T) - 2)(n - 1)}{n} \right\rfloor \geq n . \]

The result thus follows from Lemma 9.4. \hfill \Box

For every tree \( T \), let \( \text{bal}(T) \) be the maximum order of a balance subtree in \( T \). The next lemma shows how to construct a balance of large order. For each vertex \( r \) of \( T \), let \( T_r \) be the component of \( T - r \) with the maximum number of vertices.

**Lemma 9.5.** Let \( r \) be a vertex in a tree \( T \) with degree \( d \geq 3 \). Then every balance in \( T \) rooted at \( r \) has order at most \( v(T) - v(T_r) - 2 \). On the other hand, there is a balance in \( T \) rooted at \( r \) of order at least \( \frac{1}{3}(v(T) - v(T_r) - 1) \).

**Proof.** Consider a balance rooted at \( r \). The largest component of \( T - r \) is contained in at most one branch. Thus the other branch has at most \( v(T) - v(T_r) - 2 \) vertices. Hence the order of the balance is at most \( v(T) - v(T_r) - 2 \).

Now we prove the second claim. Let \( Y_1, \ldots, Y_d \) be the components of \( T - r \). Say \( v(Y_1) \geq \cdots \geq v(Y_d) \). Then \( v(T_r) = v(Y_1) \). Let \( s \) be the neighbour of \( r \) in \( Y_d \).

First suppose that \( d = 3 \). Then \( Y_1 \) and \( Y_2 \) are the branches of a balance rooted at \( r \) with support \( s \), and order \( v(Y_2) \). Now \( 2v(Y_2) \geq v(Y_2) + v(Y_3) = v(T) - v(Y_1) - 1 \). Thus the order of the balance, \( v(Y_2) \), is at least \( \frac{1}{2}(v(T) - v(Y_1) - 1) > \frac{1}{3}(v(T) - v(Y_1) - 1) \).

Now assume that \( d \geq 4 \). If \( A \) and \( B \) are a partition of \([d - 1]\), then \( \cup \{ Y_i : i \in A \} \) and \( \cup \{ Y_i : i \in B \} \) are the branches of a balance rooted at \( r \) with support \( s \). The order of the balance is
\[
\min \left\{ \sum_{i \in B} v(Y_i), \sum_{i \in A} v(Y_i) \right\} .
\]

A greedy algorithm\(^5\) gives such a partition with
\[
\min \left\{ \sum_{i \in B} v(Y_i), \sum_{i \in A} v(Y_i) \right\} \geq \frac{1}{2} \left( -v(Y_1) + \sum_{i \in [d-1]} v(Y_i) \right) = \frac{1}{2}(v(T) - v(Y_1) - v(Y_d) - 1) .
\]

Since \( Y_d \) is the smallest of \( Y_2, \ldots, Y_d \), the order of the balance is at least
\[
\frac{1}{2}(v(T) - v(Y_1) - \frac{v(T) - v(Y_1) - 1}{d - 1}) = \frac{d - 2}{2(d-1)}(v(T) - v(Y_1) - 1) \geq \frac{1}{3}(v(T) - v(Y_1) - 1) ,
\]
as desired. \hfill \Box

Which trees \( T \) have small \( \text{bal}(T) \)? First note that \( \text{bal}(P) = 0 \) for every path \( P \). A path \( P \) in a tree \( T \) is clean if every internal vertex of \( P \) has degree 2 in \( T \). Let \( p(T) \) be the maximum number of vertices in a clean path in \( T \). The hangover of \( T \), denoted by

\(^5\)Given integers \( m_1 \geq m_2 \geq \cdots \geq m_t \geq 1 \) that sum to \( m \), construct a partition of \([t]\) into sets \( A \) and \( B \) as follows. For \( i \in [t] \), let \( A_i := \sum_{j \in [i] \cap A} m_j \) and \( B_i := \sum_{j \in [i] \cap B} m_j \). Initialise \( A := \{m_1\} \) and \( B := \emptyset \). Then for \( i = 2, 3, \ldots, t \), if \( A_{i-1} \leq B_{i-1} \) then add \( i \) to \( A \); otherwise add \( i \) to \( B \). Thus \( |A_t - B_t| = m_1 \) and by induction, \( |A_t - B_t| \leq m_t \). Hence \( A_t \) and \( B_t \) are both at least \( \frac{1}{2}(m - m_1) \).
hang(T), is the minimum, taken over all clean paths P in T, of the maximum number of vertices in a component of T − E(P). We now prove that bal(T) and hang(T) are tied.

**Lemma 9.6.** For every tree T,
\[ \text{bal}(T) + 1 \leq \text{hang}(T) \leq 3 \text{bal}(T) + 1. \]

**Proof.** First we prove the lower bound. Let P be a longest clean path in T. Since every internal vertex in P has degree 2 in T, every balance in T is rooted at a vertex in one of the components of the tree obtained by deleting the internal vertices and edges of P from T. For every such vertex v, T − v has a component of at least v(T) − hang(T) − 1 vertices. That is, v(Tv) ≥ v(T) − hang(T) − 1. By Lemma 9.5, every balance rooted at v has order at least v(T) − v(Tv) − 2 ≥ hang(T) − 1. Thus \( \text{bal}(T) \leq \text{hang}(T) − 1 \).

Now we prove the upper bound. If T is a path then \( \text{bal}(T) = \text{hang}(T) = 0 \), and we are done. Now assume that T has a vertex of degree at least 3. Let r be a vertex of degree at least 3 in T such that \( v(T_r) \) is minimised. Let x be the closest vertex in \( T_r \) to r such that \( \deg(x) \neq 2 \). Let P be the path between r and x in T. Thus P is clean. If x is a leaf, then \( v(T_x) = v(T) − 1 \geq v(T_r) \). If \( \deg(x) \geq 3 \) then \( v(T_x) \geq v(T_r) \) by the choice of r. In both cases \( v(T_x) \geq v(T_r) \), which implies that r is in \( T_x \). Thus deleting the internal vertices and edges of P gives a forest with two components, one with \( v(T) − v(T_r) \) vertices, and the other with \( v(T) − v(T_x) \) vertices. Hence \( \text{hang}(T) \leq \max\{v(T) − v(T_r), v(T) − v(T_x)\} = v(T) − v(T_r) \). By Lemma 9.5, there is a balance in T rooted at r of order at least \( \frac{1}{3}(v(T) − v(T_r) − 1) \geq \frac{1}{3}(\text{hang}(T) − 1) \). Hence \( \text{bal}(T) \geq \frac{1}{3}(\text{hang}(T) − 1) \).

We now prove that if the product of sufficiently large trees has bounded Hadwiger number then both trees have bounded hangover.

**Lemma 9.7.** Fix an integer \( c \geq 1 \). Let \( T_1 \) and \( T_2 \) be trees, such that \( v(T_1) \geq 2c^2 − c + 2 \), \( v(T_2) \geq c + 1 \), and \( \eta(T_1 \square T_2) \leq c \). Then
\[ \text{hang}(T_2) \leq 3c + 1. \]

By symmetry, if in addition \( v(T_2) \geq 2c^2 − c + 2 \) then
\[ \text{hang}(T_1) \leq 3c + 1. \]

**Proof.** If \( \text{star}(T_1) \geq c \), then by Corollary 5.3, \( \eta(T_1 \square T_2) \geq \min\{v(T_2), \text{star}(T_1) + 1\} \geq c + 1 \), which contradicts the assumption. Now assume that \( \text{star}(T_1) \leq c − 1 \).

Let \( n := c + 1 \). Then \( v(T_1) \geq (c + 1)^2 + (c − 3)c + 1 + 1 = n^2 + (\text{star}(T_1) − 2)(n − 1) + 1 \). Thus Theorem 9.2 is applicable to \( T_1 \) with \( n = c + 1 \). Hence \( T_1 \) has \( c + 1 \) disjoint subtrees, each with at least \( c + 1 \) vertices.

If \( \text{bal}(T_2) \geq c + 1 \), then by Theorem 9.1, \( \eta(T_1 \square T_2) \geq \min\{c + 1, \text{bal}(T_2)\} = c + 1 \), which contradicts the assumption. Thus \( \text{bal}(T_2) \leq c \), and by Lemma 9.6, \( \text{hang}(T_2) \leq 3c + 1 \). □

We now prove a converse result to Lemma 9.7. It says that the product of two trees has small Hadwiger number whenever one of the trees is small or both trees have small hangover.
Lemma 9.8. Let $T_1$ and $T_2$ be trees, such that for some integer $c \geq 1$,

- $v(T_1) \leq c$ or $v(T_2) \leq c$, or
- $\text{hang}(T_1) \leq c$ and $\text{hang}(T_2) \leq c$.

Then $\eta(T_1 \Box T_2) \leq c'$ for some $c'$ depending only on $c$.

Proof. First suppose that $v(T_1) \leq c$. Then $\eta(T_1 \Box T_2) \leq \eta(K_c \Box T_2) = c + 1$ by Theorem 10.1 below. Similarly, if $v(T_2) \leq c$ then $\eta(T_1 \Box T_2) \leq c + 1$.

Otherwise, by assumption, $\text{hang}(T_1) \leq c$ and $\text{hang}(T_2) \leq c$. For $i \in [2]$, let $P_i$ be a clean path with $p(T_i)$ vertices in $T_i$. Now $P_1 \Box P_2$ is a planar $p(T_1) \times p(T_2)$ grid subgraph $H$ in $T_1 \Box T_2$. We now show that $T_1 \Box T_2$ can be obtained from $H$ by adding a vortex in the outerface of $H$.

Let $F$ be the set of vertices on the outerface of $H$ in clockwise order. Consider a vertex $v = (v_1, v_2) \in F$, where $v_1 \in V(P_1)$ and $v_2 \in V(P_2)$. For $i \in [2]$, let $A_i(v)$ be the component of $T_i - E(P_i)$ that contains $v_i$. As illustrated in Figure 7, define $S(v)$ to be the set $\{(a, b) : a \in A_1(v), b \in A_2(v)\}$ of vertices in $T_1 \Box T_2$. Every vertex of $G - H$ is in $S(v)$ for some vertex $v \in F$. In addition, each vertex $v \in F$ is in $S(v)$. For every edge $e$ of $T_1 \Box T_2$ where both endpoints of $e$ are in $\cup\{S(v) : v \in F\}$, the endpoints of $e$ are in one bag or in bags corresponding to consecutive vertices in $F$. For each vertex $v \in F$, if $w$ is clockwise from $v$ in $F$, then define $S'(v) := S(v) \cup S(w)$. Hence for every edge $xy$ of $T_1 \Box T_2$ where both endpoints of $e$ are in $\cup\{S(v) : v \in F\}$, the endpoints of $e$ are both in $S'(v)$ for some vertex $v \in F$. Now $|S(v)| = |A_1(v)| \cdot |A_2(v)| \leq \text{hang}(T_1)^2 \leq c^2$. Thus $\{S'(v) : v \in F\}$ is a vortex of width at most $2c^2$.

Joret and Wood [29] proved that every graph obtained from a graph embedded in a surface of bounded genus by adding a vortex of bounded width has bounded Hadwiger number. (This result was known in the graph minors community, but was never proved formally.) Thus $\eta(T_1 \Box T_2)$ is at most some constant depending only on $c$. \qed

Lemmas 9.7 and 9.8 imply the following rough structural characterisation theorem for the products of trees.

Theorem 9.9. Let $T_1$ and $T_2$ be trees (or more precisely, infinite families of trees). Then $\eta(T_1 \Box T_2)$ is bounded if and only if:

- $v(T_1)$ or $v(T_2)$ is bounded, or
- $\text{hang}(T_1)$ and $\text{hang}(T_2)$ are bounded.

10. PRODUCT OF A GENERAL GRAPH AND A COMPLETE GRAPH

This section studies the Hadwiger number of the product of a general graph and a complete graph. Miller [35] stated without proof that $\eta(T \Box K_n) = n + 1$ for every tree $T$ and integer $n \geq 2$. We now prove this claim.

Theorem 10.1. For every tree $T$ with at least one edge and integer $n \geq 1$,

$$\eta(T \Box K_n) = n + 1.$$
Proof. Since $K_2 \square K_n$ is a subgraph of $T \square K_n$, the lower bound $\eta(T \square K_n) \geq n + 1$ follows from Proposition 7.2. It remains to prove the upper bound $\eta(T \square K_n) \leq n + 1$. Let $X_1, \ldots, X_k$ be the branch sets of a complete minor in $T \square K_n$, where $k = \eta(T \square K_n)$. For each $i \in [k]$, let $T_i$ be the subtree of $T$ consisting of the edges $vw \in E(T)$ such that $(v,j)$ or $(w,j)$ is in $X_i$ for some $j \in [n]$. Since $X_i$ is connected, $T_i$ is connected. Since $X_i$ and $X_j$ are adjacent, $T_i$ and $T_j$ share an edge in common. By the Helly property of trees, there is an edge $vw$ of $T$ in every subtree $T_i$. Let $Y$ be the set of vertices $\{(v,j), (w,j) : j \in [n]\}$. Thus, by construction, every $X_i$ contains a vertex in $Y$. Since $|Y| = 2n$, if every $X_i$ has at least two vertices in $Y$, then $k \leq n$, and we are done. Now assume that some $X_i$ has only one vertex in $Y$. Say $(v,j)$ is the vertex in $X_i \cap Y$. Let $T_v$ and $T_w$ be the subtrees of $T$ obtained by deleting the edge $vw$, where $v \in V(T_v)$ and $w \in V(T_w)$. Thus $X_i$ is contained in $T_v \square K_n$. Let $Z$ be the set of neighbours of $(v,j)$ in $Y$. That is, $Z = \{(v,\ell) : \ell \in [n] \setminus \{j\}\} \cup \{(w,j)\}$. Suppose on the contrary that some branch set $X_p (p \neq i)$ has no vertex in $Z$. Then $X_p$ is contained in $T_w \square K_n$ minus the vertex $(w,j)$. Thus $X_i$ and $X_p$ are not adjacent. This contradiction proves that every branch set $X_p (p \neq i)$ has a vertex in $Z$. Since $|Z| = n$, $k \leq n + 1$, as desired. \[\square\]

Theorem 10.1 is generalised through the notion of treewidth. A tree decomposition of a graph $G$ consists of a tree $T$ and a set $\{T_x \subseteq V(G) : x \in V(T)\}$ of ‘bags’ of vertices of $G$ indexed by $T$, such that

- for each edge $vw \in E(G)$, there is some bag $T_x$ that contains both $v$ and $w$, and
- for each vertex $v \in V(G)$, the set $\{x \in V(T) : v \in T_x\}$ induces a non-empty (connected) subtree of $T$.
The *width* of the tree decomposition is $\max\{|T_x| : x \in V(T)\} - 1$. The *treewidth* of $G$ is the minimum width of a tree decomposition of $G$. For example, $G$ has treewidth 1 if and only if $G$ is a forest.

**Theorem 10.2.** For every graph $G$ and integer $n \geq 1$,

$$\eta(G \Box K_n) \leq \text{tw}(G \Box K_n) + 1 \leq n(\text{tw}(G) + 1).$$

Moreover, for all integers $k \geq 2$ and $n \geq 2$ there is a graph $G$ with $\text{tw}(G) = k$ and

$$\eta(G \Box K_n) = \text{tw}(G \Box K_n) + 1 = n(k + 1).$$

**Proof.** First we prove the upper bound. Let $(T, \{T_x \subseteq V(G) : x \in V(T)\})$ be a tree decomposition of $G$ with at most $\text{tw}(G) + 1$ vertices in each bag. Replace each bag $T_x$ by $\{(v, i) : v \in T_x, i \in [n]\}$. We obtain a tree decomposition of $G \Box K_n$ with at most $n(\text{tw}(G) + 1)$ vertices in each bag. Thus $\text{tw}(G \Box K_n) \leq n(\text{tw}(G) + 1) - 1$. Every graph $H$ satisfies $\eta(H) \leq \text{tw}(H) + 1$. The result follows.

Now we prove the lower bound. Let $G$ be the graph with vertex set

$$V(G) := \{v_1, \ldots, v_{k+1}\} \cup \{x_{i,j,p} : i, j \in [k+1], p \in [n]\},$$

where $\{v_1, \ldots, v_{k+1}\}$ is a clique, and each $x_{i,j,p}$ is adjacent to $v_i$ and $v_j$. A tree decomposition $T$ of $G$ is constructed as follows. Let $T_r := \{v_1, \ldots, v_{k+1}\}$, and for all $i, j \in [k+1]$ and $p \in [n]$, let $T_{i,j,p} := \{x_{i,j,p}, v_i, v_j\}$, where $T_r$ is adjacent to every $T_{i,j,p}$ and there are no other edges in $T$. Thus $T$ is a star with $n(k+1)^2$ leaves, and $(T, \{T_x : x \in V(T)\})$ is a tree decomposition of $G$ with at most $k + 1$ vertices in each bag. Thus $\text{tw}(G) \leq k$. Since $G$ contains a clique of $k + 1$ vertices, $\text{tw}(G) = k$.

Now consider $G \Box K_n$. For $i, j \in [k+1]$ and $p \in [n]$, let $A(i, j, p)$ be the subgraph of $G \Box K_n$ induced by $\{(x_{i,j,p}, q) : q \in [n]\}$. Thus $A(i, j, p)$ is a copy of $K_n$. For $i \in [k+1]$ and $p \in [n]$, let $X(i, p)$ be the subgraph induced by $\cup\{A(i, j, p) : j \in [k+1]\}$ plus the vertex $\{v_i, p\}$. We claim that the $X(i, p)$ are the branch set of clique minor in $G \Box K_n$. First we prove that each $X(i, p)$ is connected. For all $j \in [k + 1]$, $v_i$ is adjacent to $x_{i,j,p}$ in $G$. Thus $(v_i, p)$ is adjacent to $(x_{i,j,p}, p)$, which is in $A(i, j, p) \subset X(i, p)$. Thus $X(i, p)$ consists of $k + 1$ copies of $K_n$ plus one vertex adjacent to each copy. In particular, $X(i, p)$ is connected. Now consider distinct subgraphs $X(i, p)$ and $X(j, q)$. The first coordinate of every vertex in $X(i, p)$ is either $v_{i,p}$ or $x_{i',p}$ for some $i' \in [k + 1]$. Thus $X(i, p)$ and $X(j, q)$ are disjoint. Now the vertex $x_{i,j,p}$ is adjacent to the vertex $v_{j,q}$ in $G$. Thus the vertex $(x_{i,j,p}, q)$, which is in $A(i, j, p) \subset X(i, p)$, is adjacent to the vertex $(v_{j,q}, q)$, which is in $X(j, q)$. Thus $X(i, p)$ and $X(j, q)$ are adjacent. Hence the $X(i, p)$ are the branch set of clique minor in $G \Box K_n$, and $\eta(G \Box K_n) \geq n(k + 1)$. We have equality because of the above upper bound.

We now set out to prove a lower bound on $\eta(G \Box K_n)$ in terms of the treewidth of $G$. We start by considering the case $n = 2$, which is of particular importance in Section 12 below. Robertson and Seymour [40] proved that every graph with large treewidth has a

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6This upper bound even holds for strong products.
large grid minor. The following explicit bound was obtained by Diestel et al. \cite{12}; also see \cite{11}, Theorem 12.4.4.

**Lemma 10.3** (\cite{12}). For all integers \( k, m \geq 1 \) every graph with tree-width at least \( k^{4m^2(k+2)} \) contains \( P_k \square P_k \) or \( K_m \) as a minor. In particular, every graph with tree-width at least \( k^{4k^4(k+2)} \) contains a \( P_k \square P_k \)-minor.

In what follows all logarithms are binary.

**Lemma 10.4.** For every graph \( G \) with at least one edge,

\[
\eta(G \square K_2) > \left( \frac{4}{7} \log \text{tw}(G) \right)^{1/4}.
\]

**Proof.** Let \( \ell \) be the real-valued solution to \( \text{tw}(G) = \ell^{4(\ell+1)^2} \). Thus \( \ell \geq 1 \), and

\[
\log \text{tw}(G) = 4(\ell + 1)^3(\ell) < 4(\ell + 1)^4.
\]

That is, \( \left( \frac{4}{7} \log \text{tw}(G) \right)^{1/4} < \ell + 1 \). Let \( k := \lceil \ell \rceil \). Thus \( k \geq 1 \) and \( \text{tw}(G) \geq k^{4(k+1)^3} \). Hence Lemma 10.3 is applicable with \( m = k + 1 \). Thus \( G \) contains \( P_k \square P_k \) or \( K_{k+1} \) as a minor. If \( G \) contains a \( P_k \square P_k \)-minor, then \( G \square K_2 \) contains a \( K_{k+2} \)-minor by (1). Otherwise \( G \) contains a \( K_{k+1} \) minor, and by Proposition 7.2 \( G \square K_2 \) contains a \( K_{k+2} \)-minor. In both cases

\[
\eta(G \square K_2) \geq k + 2 > \ell + 1 > \left( \frac{4}{7} \log \text{tw}(G) \right)^{1/4},
\]

as desired. \( \square \)

Lemma 10.4 and Theorem 10.2 imply that \( \eta(G \square K_2) \) is tied to the treewidth of \( G \). In particular,

\[
(4) \quad \left( \frac{4}{7} \log \text{tw}(G) \right)^{1/4} < \eta(G \square K_2) \leq 2 \text{tw}(G) + 2.
\]

We now extend Lemma 10.4 for general complete graphs.

**Lemma 10.5.** For every graph \( G \) with at least one edge and every integer \( n \geq 1 \),

\[
\eta(G \square K_n) > \frac{n}{16} \left( \frac{4}{7} \log \text{tw}(G) \right)^{1/6}.
\]

**Proof.** Let \( \ell \) be the real-valued solution to \( \text{tw}(G) = \ell^{4(\ell+2)^3} \). Thus \( \ell \geq 1 \). Thus

\[
\log \text{tw}(G) = 4\ell^4(\ell+2)(\ell) \leq 12\ell^6.
\]

That is, \( \left( \frac{4}{7} \log \text{tw}(G) \right)^{1/6} \leq \ell \). Let \( k := \lceil \ell \rceil \). Thus \( k \geq \ell \geq 1 \) and \( \text{tw}(G) \geq k^{4k^4(k+2)} \). By Lemma 10.3 \( G \) contains a \( P_k \square P_k \)-minor. By Lemma 10.6 below,

\[
\eta(G \square K_n) \geq \frac{n}{16} (k + 1) > \frac{n}{16} \ell \geq \frac{n}{16} \left( \frac{4}{7} \log \text{tw}(G) \right)^{1/6},
\]

as desired. \( \square \)

Lemma 10.5 and Theorem 10.2 imply that \( \eta(G \square K_n)/n \) is tied to the treewidth of \( G \). In particular,

\[
\left( \frac{4}{7} \log \text{tw}(G) \right)^{1/6} \leq \frac{\eta(G \square K_n)}{n} \leq \text{tw}(G) + 1.
\]

It remains to prove Lemma 10.6.
Lemma 10.6. For all integers $n \geq 1$ and $k \geq 1$

$$(k+1)\left\lceil \frac{n}{2} \right\rceil \leq \eta(P_k \Box P_k \Box K_n) < k(n + \frac{1}{k}) + 3.$$ 

Proof. Since $P_k \Box P_k \Box K_n$ has $k^2n$ vertices and maximum degree $n + 3$, Lemma 2.2(c) implies the upper bound, $\eta(P_k \Box P_k \Box K_n) \leq \sqrt{(n + 1)k^2n} + 3 < k(n + \frac{1}{k}) + 3$.

Now we prove the lower bound. Let $p := \lfloor \frac{n}{2} \rfloor$. Each vertex is described by a triple $(x, y, r)$ where $x, y \in [k]$ and $r \in [n]$. Distinct vertices $(x, y, r)$ and $(x', y', r')$ are adjacent if and only if $x = x'$ and $y = y'$, or $x = x'$ and $|y - y'| = 1$ and $r = r'$, or $y = y'$ and $|x - x'| = 1$ and $r = r'$.

For $r \in [p]$, let $T_{0,r}$ be the subgraph induced by $\{(1, y, 2r - 1) : y \in [k]\}$, and let $T_{1,r}$ be the subgraph induced by $\{(x, 1, 2r) : x \in [k]\}$. For $i \in [2, k]$ and $r \in [p]$, let $T_{i,r}$ be the subgraph of $P_k \Box P_k \Box K_n$ induced by $\{(i, y, 2r - 1) : y \in [k]\} \cup \{(x, i, 2r) : x \in [k]\}$.

For all $r \in [p]$, both $T_{0,r}$ and $T_{1,r}$ are paths, and for $i \in [2, k]$, $T_{i,r}$ consists of two adjacent paths. In particular, each $T_{i,r}$ is connected. Observe that each pair of distinct subgraphs $T_{i,r}$ and $T_{j,s}$ are disjoint.

There is an edge from $(1, 1, 2r - 1)$ in $T_{0,r}$ to $(1, 1, 2s)$ in $T_{1,s}$. For all $i \in [2, k]$, there is an edge from $(i, 1, 2r - 1)$ in $T_{1,r}$ to $(1, i, 2s - 1)$ in $T_{0,s}$, and for all $i, j \in [2, k]$, there is an edge from $(j, i, 2r)$ in $T_{i,r}$ to $(j, i, 2s - 1)$ in $T_{j,s}$.

Hence the $T_{i,j}$ are the branch sets of a $K_{(k+1)m}$-minor. Therefore $\eta(P_k \Box P_k \Box K_n) \geq (k+1)p = (k+1)\left\lceil \frac{n}{2} \right\rceil$. $\square$

11. Rough Structural Characterisation Theorem

In this section we give a rough structural characterisation of pairs of graphs whose product has bounded Hadwiger number. The proof is based heavily on the corresponding result for trees in Section 9. Thus our first task is to extend a number of definitions for trees to general graphs.

For a connected graph $G$, let $bal(G)$ be the maximum order of a balance subgraph in $G$. A path $P$ in $G$ is semi-clean if every internal vertex of $P$ has degree 2 in $G$. Let $p'(G)$ be the maximum number of vertices in a semi-clean path in $G$. A path $P$ is clean if it is semi-clean, and every edge of $P$ is a cut in $G$. Let $p(G)$ be the maximum number of vertices in a clean path in $G$. Note that $p(G) \geq 1$ since a single vertex is a clean path. In fact, if $G$ is 2-connected, then the only clean paths are single vertices, and $p(G) = 1$. On the other hand, since every edge in a tree $T$ is a cut, our two definitions of a clean path are equivalent for trees, and $p'(T) = p(T)$.

The hangover of a connected graph $G$, denoted by $hang(G)$, is defined as follows. If $G$ is a path or a cycle then $hang(G) := 0$. Otherwise, $hang(G)$ is the minimum, taken over all clean paths $P$ in $G$, of the maximum number of vertices in a component of $G - E(P)$.

First note the following trivial relationship between $hang(G)$ and $p(G)$:

$$\frac{1}{2}(v(G) - p(G) + 2) \leq hang(G) \leq v(G) - p(G) + 1. \quad (5)$$
To prove a relationship between \( \text{bal}(G) \) and \( \text{hang}(G) \) below, we reduce the proof to the case of trees using the following lemma.

**Lemma 11.1.** Every connected graph \( G \) has a spanning tree \( T \) such that \( p(T) \leq p'(G) + 6 \).

**Proof.** Define a leaf-neighbour in a tree to be a vertex of degree 2 that is adjacent to a leaf (a vertex of degree 1).

Choose a spanning tree \( T \) of \( G \) that firstly maximises the number of leaves in \( T \), and secondly maximises the number of leaf-neighbours in \( T \).

If \( p(T) \leq 7 \) then the claim is vacuous since \( p'(G) \geq 1 \). Now assume that \( p(T) \geq 8 \). Let \((v_1, \ldots, v_k)\) be a clean path in \( T \) with \( k = p(T) \). Below we prove that \( \deg_G(v_i) = 2 \) for each \( i \in [4, k - 3] \). This shows that the path \((v_4, \ldots, v_{k-3})\) is semi-clean in \( G \), implying \( p'(G) \geq p(T) - 6 \), as desired.

Suppose on the contrary that \( \deg_G(v_i) \geq 3 \) for some \( i \in [4, k - 3] \). Let \( w \) be a neighbour of \( v_i \) in \( G \) besides \( v_{i-1} \) and \( v_{i+1} \). Without loss of generality, the path between \( v_i \) and \( w \) in \( T \) includes \( v_{i+1} \).

**Case 1.** \( \deg_T(w) \geq 2 \): Let \( T' \) be the spanning tree of \( G \) obtained from \( T \) by deleting the edge \( v_i v_{i+1} \) and adding the edge \( v_i w \). Now \( \deg_T(v_{i+1}) = 2 \) (since \( i + 1 \leq k - 1 \)). Thus \( v_{i+1} \) becomes a leaf in \( T' \). Since \( \deg_T(v_i) = \deg_T(v_i) = 2 \), \( v_i \) is a leaf in neither \( T \) nor \( T' \). Since \( \deg_T(w) \geq 2 \) and \( \deg_T(w) \geq 3 \), \( w \) is a leaf in neither \( T \) nor \( T' \). The degree of every other vertex is unchanged. Hence \( T' \) has one more leaf than \( T \). This contradicts the choice of \( T \).

Now assume that \( \deg_T(w) = 1 \). Let \( x \) be the neighbour of \( w \) in \( T \).

**Case 2.** \( \deg_T(w) = 1 \) and \( \deg_T(x) = 2 \): Let \( T' \) be the spanning tree of \( G \) obtained from \( T \) by deleting the edge \( wx \) and adding the edge \( v_i w \). Since \( \deg_T(v_i) = 2 \) and \( \deg_T(v_i) = 3 \), \( v_i \) is a leaf in neither \( T \) nor \( T' \). Since \( \deg_T(w) = \deg_T(w) = 1 \), \( w \) is a leaf in both \( T \) and \( T' \). Since \( \deg_T(x) = 2 \) and \( \deg_T(x) = 1 \), \( x \) becomes a leaf in \( T' \). The degree of every other vertex is unchanged. Hence \( T' \) has one more leaf than \( T \). This contradicts the choice of \( T \).

**Case 3.** \( \deg_T(w) = 1 \) and \( \deg_T(x) \geq 3 \): Let \( T' \) be the spanning tree of \( G \) obtained from \( T \) by deleting the edge \( v_i v_{i+1} \) and adding the edge \( v_i w \). Since \( \deg_T(v_i) = \deg_T(v_i) = 2 \), \( v_i \) is a leaf in neither \( T \) nor \( T' \). Now \( \deg_T(v_{i+1}) = 2 \) (since \( i + 1 \leq k - 1 \)). Thus \( v_{i+1} \) is a leaf in \( T' \) but not in \( T \). Since \( \deg_T(w) = 1 \) and \( \deg_T(w) = 2 \), \( w \) is a leaf in \( T \) but not in \( T' \). The degree of every other vertex is unchanged. Hence \( T' \) has the same number of leaves as \( T \).

Suppose, for the sake of contradiction, that there is a leaf-neighbour \( p \) in \( T \) that is not a leaf-neighbour in \( T' \). Since \( v_{i+1} \) and \( w \) are the only vertices with different degrees in \( T \) and \( T' \), \( p \) is either \( v_{i+1} \) or \( w \), or \( p \) is a neighbour of \( v_{i+1} \) or \( w \) in \( T \) or \( T' \). That is, \( p \in \{v_{i+1}, w, x, v_i, v_{i+2}\} \). Now \( \deg_T(p) = 2 \) since \( p \) is a leaf-neighbour in \( T \). Thus \( p \neq w \) and \( p \neq x \). Every neighbour of \( v_i \) and \( v_{i+1} \) in \( T \) has degree 2 in \( T \) (since \( i - 1 \geq 2 \) and \( i + 2 \leq k - 1 \)). Thus \( p \neq v_i \) and \( p \neq v_{i+1} \). Finally, \( p \neq v_{i+2} \) since the neighbours of \( v_{i+2} \) in \( T \), namely \( v_{i+1} \) and \( v_{i+3} \), are both not leaves (since there is a path in \( T \) from \( v_{i+3} \) to
x that avoids \(v_{i+2}\). This contradiction proves that every leaf-neighbour in \(T\) is also a leaf-neighbour in \(T'\).

Now consider the vertex \(v_{i+2}\). In both \(T\) and \(T'\), the only neighbours of \(v_{i+2}\) are \(v_{i+1}\) and \(v_{i+3}\) (since \(i + 2 \leq k - 1\)). Both \(v_{i+1}\) and \(v_{i+3}\) have degree 2 in \(T\), but \(v_{i+1}\) is a leaf in \(T'\). Thus \(v_{i+2}\) is a leaf-neighbour in \(T'\), but not in \(T\).

Hence \(T'\) has more leaf-neighbours than \(T\). This contradicts the choice of \(T\), and completes the proof. \(\square\)

Lemma 9.6 proves that \(\text{bal}(T)\) and \(\text{hang}(T)\) are tied for trees. We now prove an analogous result for general graphs.

**Lemma 11.2.** For every connected graph \(G\),

\[
\text{hang}(G) \leq 8 \text{bal}(G) + 9.
\]

**Proof.** If \(G\) is a path or cycle, then \(\text{hang}(G) = 0\) and the result is vacuous. Now assume that \(G\) is neither a path nor a cycle. By Lemma [11.1], \(G\) has a spanning tree \(T\) such that \(p(T) \leq p'(G) + 6\). By Lemma [9.6]

\[
\text{bal}(G) \geq \text{bal}(T) \geq \frac{1}{3}(\text{hang}(T) - 1).
\]

By the lower bound in (5),

\[
\text{bal}(G) \geq \frac{1}{3}(\frac{1}{2}(\nu(T) - p(T) + 2) - 1) = \frac{1}{6}(\nu(G) - p(T)) \geq \frac{1}{6}(\nu(G) - p'(G) - 6).
\]

Thus we are done if \(\frac{1}{6}(\nu(G) - p'(G) - 6) \geq \frac{1}{8}(\text{hang}(G) - 9)\). Now assume that

\[
\frac{1}{6}(\nu(G) - p'(G) - 6) \leq \frac{1}{8}(\text{hang}(G) - 9).
\]

By the upper bound in (5),

\[
\frac{1}{6}(\nu(G) - p'(G) - 6) \leq \frac{1}{8}(\nu(G) - p(G) + 1 - 9).
\]

That is,

\[
\nu(G) + 3p(G) \leq 4p'(G).
\]

If \(p(G) \geq p'(G)\), then \(\nu(G) \leq p'(G)\), which implies that \(G\) is a path. Now assume that \(p(G) \leq p'(G) - 1\). Thus there is a non-clean semi-clean path \(P\) in \(G\) of length \(p'(G)\). Since \(P\) is not clean and \(G\) is connected and not a cycle, there is a cycle \(C\) in \(G\) with at least \(p'(G)\) vertices, such that one vertex \(r\) in \(C\) is adjacent to a vertex \(s\) not in \(C\). It follows that \(G\) has a balance rooted at \(r\) with support \(s\), and with order at least \(\lfloor \frac{1}{2}(p'(G) - 1) \rfloor\).

Thus \(\text{bal}(G) \geq \frac{1}{2}(p'(G) - 2)\). That is, \(8 \text{bal}(G) + 8 \geq 4p'(G)\). By (5),

\[
8 \text{bal}(G) + 8 \geq \nu(G) + 3p(G) \geq \nu(G) \geq \text{hang}(G),
\]

as desired. \(\square\)

We now prove an analogue of Lemma 9.7 for general graphs.
Lemma 11.3. Fix an integer $c \geq 1$. Let $G$ and $H$ be graphs, such that $v(G) \geq 2c^2 - c + 2$, $v(H) \geq c + 1$, and $\eta(G \Box H) \leq c$. Then
\[
\text{hang}(H) \leq 8c + 9.
\]

By symmetry, if in addition $v(H) \geq 2c^2 - c + 2$ then
\[
\text{hang}(G) \leq 8c + 9.
\]

Proof. If $\text{star}(G) \geq c$, then by Corollary 5.3 $\eta(G \Box H) \geq \min\{v(H), \text{star}(G) + 1\} \geq c + 1$, which contradicts the assumption. Now assume that $\text{star}(G) \leq c - 1$.

Let $T$ be a spanning tree of $G$. Let $n := c + 1$. Then
\[
v(T) = v(G) \geq (c + 1)^2 + (c - 3)c + 1 \geq n^2 + (\text{star}(G) - 2)(n - 1) + 1 \\
\geq n^2 + (\text{star}(T) - 2)(n - 1) + 1.
\]

Thus Theorem 9.2 is applicable to $T$ with $n = c + 1$. Hence $T$ has $c + 1$ disjoint subtrees, each with at least $c + 1$ vertices.

If $\text{bal}(H) \geq c + 1$, then by Theorem 9.1 $\eta(G \Box H) \geq \min\{c + 1, \text{bal}(H)\} = c + 1$, which contradicts the assumption. Thus $\text{bal}(H) \leq c$. Hence, by Lemma 11.2 $\text{hang}(H) \leq 8c + 9$. \hfill \Box

We now prove that the product of graphs with bounded hangover have a specific structure.

Lemma 11.4. Fix an integer $c \geq 1$. For all graphs $G$ and $H$, if $\text{hang}(G) \leq c$ and $\text{hang}(H) \leq c$, then $G \Box H$ is one of the following graphs:

- a planar grid (the product of two paths) with a vortex of width at most $2c^2$ in the outerface,
- a cylindrical grid (the product of a path and a cycle) with a vortex of width at most $2c$ in each of the two ‘big’ faces, or
- a toroidal grid (the product of two cycles).

Proof. If $G$ and $H$ are cycles then $G \Box H$ is a toroidal grid. If neither $G$ nor $H$ are cycles then by the same argument used in the proof of Theorem 9.9 $G \Box H$ is obtained from a planar $p(G) \times p(H)$ grid by adding a vortex in the outerface with width at most $2c^2$. If $G$ is a cycle and $H$ is not a cycle, then by a similar argument used in the proof of Theorem 9.9 $G \Box H$ is obtained from a cylindrical $v(C_n) \times p(H)$ grid by adding a vortex in each of the two ‘big’ faces with width at most $2c$. \hfill \Box

Lemmas 11.3 and 11.4 imply the following characterisation of large graphs with bounded Hadwiger number that was described in Section 1.

Theorem 11.5. Fix an integer $c \geq 1$. For all graphs $G$ and $H$ with $v(G) \geq 2c^2 - c + 2$ and $v(H) \geq 2c^2 - c + 2$, if $\eta(G \Box H) \leq c$ then $G \Box H$ is one of the following graphs:

- a planar grid (the product of two paths) with a vortex of width at most $2(8c + 9)^2$ in the outerface,
• a cylindrical grid (the product of a path and a cycle) with a vortex of width at most $16c + 18$ in each of the two 'big' faces, or
• a toroidal grid (the product of two cycles).

We now prove the first part of our rough structural characterisation of graph products with bounded Hadwiger number.

**Lemma 11.6.** Let $G$ and $H$ be connected graphs, each with at least one edge, such that $\eta(G \square H) \leq c$ for some integer $c$. Then for some integers $c_1, c_2, c_3$ depending only on $c$:

- $\tw(G) \leq c_1$ and $\nu(H) \leq c_2$, or
- $\tw(H) \leq c_1$ and $\nu(G) \leq c_2$, or
- $\hang(G) \leq c_3$ and $\hang(H) \leq c_3$.

**Proof.** Let $c_1 := 2^{4c^4}$, $c_2 := 2c^2 - c + 1$, and $c_3 := 8c + 9$.

First suppose that $\tw(G) > c_1$ or $\tw(H) > c_1$. Without loss of generality, $\tw(G) > c_1$. Then by Lemma 10.4, $\eta(G \square H) \geq \eta(G \square K_2) > (\frac{4}{3} \log c_1)^{1/4} = c$, which is a contradiction. Now assume that $\tw(G) \leq c_1$ and $\tw(H) \leq c_1$.

Thus, if $\nu(H) \leq c_2$ or $\nu(G) \leq c_2$, then the first or second condition is satisfied, and we are done. Now assume that $\nu(H) > c_2$ and $\nu(G) > c_2$. By Lemma 11.3, $\hang(G)$ and $\hang(H)$ are both at most $8c + 9 = c_3$, as desired. □

Now we prove the converse of Lemma 11.6.

**Lemma 11.7.** Let $G$ and $H$ be connected graphs, each with at least one edge, such that for some integers $c_1, c_2, c_3$,

- $\tw(G) \leq c_1$ and $\nu(H) \leq c_2$, or
- $\tw(H) \leq c_1$ and $\nu(G) \leq c_2$, or
- $\hang(G) \leq c_3$ and $\hang(H) \leq c_3$.

Then $\eta(G \square H) \leq c$ where $c$ depends only on $c_1, c_2, c_3$.

**Proof.** Suppose that $\tw(G) \leq c_1$ and $\nu(H) \leq c_2$. Then Theorem 10.2 implies that

$$\eta(G \square H) \leq \eta(G \square K_{c_2}) \leq c_2(\tw(G) + 1) \leq c_2(c_1 + 1),$$

and we are done. Similarly, if $\tw(H) \leq c_1$ and $\nu(G) \leq c_2$, then $\eta(G \square H) \leq c_2(c_1 + 1)$, and we are done. Otherwise, $\hang(G) \leq c_3$ and $\hang(H) \leq c_3$. By Lemma 11.4, $G \square H$ is either a toroidal grid (which has no $K_8$ minor), or $G \square H$ is a planar graph plus vortices of width at most $2c^2$ in one or two of the faces. Robertson and Seymour [41] proved that every graph obtained from a graph embedded in a surface of bounded genus by adding vortices of bounded width has bounded Hadwiger number. Thus $\eta(G \square H)$ is at most some constant depending only on $c$. □

Lemmas 11.6 and 11.7 imply the following rough structural characterisation of graph products with bounded Hadwiger number.

**Theorem 11.8.** The function $\eta(G \square H)$ is tied to

$$\min \{ \max \{\tw(G), \nu(H)\}, \max \{\nu(G), \tw(H)\}, \max \{\hang(G), \hang(H)\} \}.$$
Theorem 11.8 can be informally stated as: \( \eta(G \square H) \) is bounded if and only if:
- \( \text{tw}(G) \) and \( \nu(H) \) is bounded, or
- \( \nu(G) \) and \( \text{tw}(H) \) is bounded, or
- \( \text{hang}(G) \) and \( \text{hang}(H) \) are bounded.

12. ON HADWIGER’S CONJECTURE FOR CARTESIAN PRODUCTS

In 1943, Hadwiger [20] made the following conjecture.

**Hadwiger’s Conjecture.** For every graph \( G \),
\[
\chi(G) \leq \eta(G).
\]

This conjecture is widely considered to be one of the most significant open problems in graph theory; see the survey by Toft [50]. Yet it is unknown whether Hadwiger’s conjecture holds for all non-trivial products. (We say \( G \square H \) is non-trivial if both \( G \) and \( H \) are both connected and have at least one edge.) The chromatic number of a product is well understood. In particular, Sabidussi [43] proved that
\[
\chi(G \square H) = \max\{\chi(G), \chi(H)\}.
\]
Thus Hadwiger’s Conjecture for products asserts that
\[
\max\{\chi(G), \chi(H)\} \leq \eta(G \square H).
\]

Hadwiger’s Conjecture is known to hold for various classes of products. For example, Chandran and Sivadasan [7] proved that the product of sufficiently many graphs (relative to their maximum chromatic number) satisfies Hadwiger’s Conjecture. The best bounds are by Chandran and Raju [5, 38], who proved that for some constant \( c \), Hadwiger’s Conjecture holds for the non-trivial product \( G_1 \square G_2 \square \cdots \square G_d \) whenever
\[
\max_i \chi(G_i) \leq 2^{2(d-c)/2}.
\]

In a different direction, Chandran and Raju [5, 38] proved that if \( \chi(G) \geq \chi(H) \) and \( \chi(H) \) is not too small relative to \( \chi(G) \), then \( G \square H \) satisfies Hadwiger’s Conjecture. In particular, there is a constant \( c \), such that if \( \chi(G) \geq \chi(H) \geq c \log^{3/2} \chi(G) \) then \( G \square H \) satisfies Hadwiger’s Conjecture. Similarly, they also implicitly proved that
\[
\min\{\chi(G), \chi(H)\} \leq \eta(G \square H),
\]
and concluded that if \( \chi(G) = \chi(H) \) then Hadwiger’s Conjecture holds for \( G \square H \). We make the following small improvement to this result.

**Lemma 12.1.** For all connected graphs \( G \) and \( H \), both with at least one edge,
\[
\min\{\chi(G), \chi(H)\} \leq \eta(G \square H) - 1.
\]
Moreover, if \( G \neq K_2 \) and \( H \neq K_2 \) then
\[
\min\{\chi(G), \chi(H)\} \leq \eta(G \square H) - 2.
\]

7 The chromatic number of a graph \( G \), denoted by \( \chi(G) \), is the minimum integer \( k \) such that each vertex of \( G \) can be assigned one of \( k \) colours such that adjacent vertices receive distinct colours.
Proof. We have $\chi(G) \leq \Delta(G) + 1$ and $\chi(H) \leq \Delta(H) + 1$. Thus by Corollary 5.3

$$\min\{\chi(G), \chi(H)\} \leq \min\{\Delta(G), \Delta(H)\} + 1 \leq \eta(G \Box H) - 1.$$  

Now assume that $G \neq K_2$ and $H \neq K_2$.

Case 1. $G \in \{C_n, K_n\}$ and $\eta(H) = 2$ for some $n \geq 3$: Then $H$ is a tree and $\eta(H) = 2$. On the other hand, $\eta(G \Box H) \geq \eta(K_3 \Box K_2) = 4$ by Proposition 7.2.

Case 2. $G = K_n$ and $\eta(H) \geq 3$ for some $n \geq 3$: Then $\chi(G) \leq \Delta(G) \leq \Delta(H) + 2$. On the other hand, $\eta(G \Box H) \geq \eta(K_3 \Box K_2) = 4$ by Proposition 7.3.

Case 3. $G = C_n$ and $\eta(H) \geq 3$ for some $n \geq 3$: Then $\min\{\chi(G), \chi(H)\} \leq 3$, and $\eta(G \Box H) \geq \eta(C_3 \Box C_3) = 5$ by a result of Archdeacon et al. [3]. (In fact, Archdeacon et al. [3] determined $\eta(C_n \Box C_m)$ for all values of $n$ and $m$, as described in Table 2. Miller [35] had previously stated without proof that $\eta(C_n \Box K_2) = 4$ for all $n \geq 3$.)

Case 4. Both $G$ and $H$ are neither complete graphs nor cycles. Then by Brooks’ Theorem 3, $\chi(G) \leq \Delta(G)$ and $\chi(H) \leq \Delta(H)$. Thus $\min\{\chi(G), \chi(H)\} \leq \min\{\Delta(G), \Delta(H)\}$.

By Corollary 5.3 $\eta(G \Box H) \geq \min\{\Delta(G), \Delta(H)\} + 2$. Thus $\min\{\chi(G), \chi(H)\} \leq \eta(G \Box H) - 2$.

\[\square\]

Table 2. The Hadwiger number of $C_n \Box C_m$ where $C_2 = K_2$; see [1].

|       | $n = 2$ | $n = 3$ | $n = 4$ | $n = 5$ | $n \geq 6$ |
|-------|---------|---------|---------|---------|-----------|
| $m = 2$ | 3       | 4       | 4       | 4       | 4         |
| $m = 3$ | 4       | 5       | 5       | 5       | 6         |
| $m = 4$ | 4       | 5       | 6       | 6       | 7         |
| $m = 5$ | 4       | 5       | 6       | 7       | 7         |
| $m \geq 6$ | 4       | 6       | 7       | 7       | 7         |

**Theorem 12.2.** Hadwiger’s Conjecture holds for a non-trivial product $G \Box H$ whenever $|\chi(G) - \chi(H)| \leq 2$. Moreover, if $|\chi(G) - \chi(H)| \leq 1$ then $\chi(G \Box H) \leq \eta(G \Box H) - 1$.

Proof. Without loss of generality, $\chi(G) - 2 \leq \chi(H) \leq \chi(G)$. Thus, by Sabidussi’s Theorem [43], it suffices to prove that $\chi(G) \leq \eta(G \Box H)$. If $H = K_2$ then $\chi(G) \leq 4$ by assumption. Hadwiger [20] and Dirac [13] independently proved Hadwiger’s Conjecture whenever $\chi(G) \leq 4$. Thus $\eta(G \Box H) \geq \eta(G) + 1 \geq \chi(G) + 1$, as desired. This proves the ‘moreover’ claim in this case. Now assume that $H \neq K_2$. Thus by Lemma 12.1 $\chi(G) \leq \chi(H) + 2 \leq \eta(G \Box H)$, as desired. And if $\chi(G) \leq \chi(H) + 1$ then $\chi(G) \leq \eta(G \Box H) - 1$, as desired. $\square$

The following two theorems establish Hadwiger’s Conjecture for new classes of products. The first says that products satisfy Hadwiger’s Conjecture whenever one graph has large treewidth relative to its chromatic number.

**Theorem 12.3.** Hadwiger’s Conjecture is satisfied for the non-trivial product $G \Box H$ whenever $\chi(G) \geq \chi(H)$ and $G$ has treewidth $\text{tw}(G) \geq 2^{4\chi(G)}$.
**Proof.** Since $H$ has at least one edge, $\eta(G \square H) \geq \eta(G \square K_2)$. By Lemma \[10.4\], $\eta(G \square H) > \left(\frac{1}{4} \log tw(G)\right)^{1/4}$, which is at least $\chi(G)$ by assumption. Hence $\eta(G \square H) > \chi(G) = \chi(G \square H)$ by Sabidussi’s Theorem \[43\]. That is, $G \square H$ satisfies Hadwiger’s Conjecture. \[\square\]

We now show that products satisfy (a slightly better bound than) Hadwiger’s Conjecture whenever the graph with smaller chromatic number is relatively large.

**Theorem 12.4.** Let $G$ and $H$ be connected graphs with $v(H) - 1 \geq \chi(G) \geq \chi(H)$. Then

$$\chi(G \square H) \leq \eta(G \square H) - \chi(G) - 1.$$  

**Proof.** By Sabidussi’s Theorem \[43\] it suffices to prove that $\eta(G \square H) \geq \chi(G) + 1$.

**Case 1.** $G = K_n$ for some $n \geq 3$: Then by Proposition \[7.2\]

$$\eta(G \square H) \geq \eta(K_n \square K_2) = n + 1 = \chi(G) + 1 = \chi(G \square H) + 1.$$  

**Case 2.** $G = C_n$ for some $n \geq 3$: Then by Proposition \[7.2\]

$$\eta(G \square H) \geq \eta(K_3 \square K_2) = 4 \geq \chi(G) + 1 = \chi(G \square H) + 1.$$  

**Case 3.** $G$ is neither a complete graph nor a cycle: Then by Brooks’ Theorem \[3\] and Corollary \[5.3\]

$$\eta(G \square H) \geq \min\{v(H), \Delta(G) + 1\} \geq \min\{v(H), \chi(G) + 1\} = \chi(G) + 1 = \chi(G \square H) + 1,$$

as desired. \[\square\]

Note that Theorem \[12.4\] is best possible, since Theorem \[10.1\] implies that for $G = K_n$ and $H$ any tree (no matter how big), $\chi(G \square H) = n = \eta(G \square H) - 1$.

Theorems \[12.2\] and \[12.4\] both prove (under certain assumptions) that $\chi(G \square H) \leq \eta(G \square H) - 1$, which is stronger than Hadwiger’s Conjecture for general graphs. This should not be a great surprise, since if Hadwiger’s Conjecture holds for all graphs, then the same improved result holds for all non-trivial products $G \square H$ with $\chi(G) \geq \chi(H)$:

$$\chi(G \square H) = \chi(G) \leq \eta(G) \leq \eta(G \square K_2) - 1 \leq \eta(G \square H) - 1.$$  

Whether Hadwiger’s Conjecture holds for all non-trivial products reduces to the following particular case.

**Theorem 12.5.** Let $G$ be a graph. Then Hadwiger’s Conjecture holds for every non-trivial product $G \square H$ with $\chi(G) \geq \chi(H)$ if and only if Hadwiger’s Conjecture holds for $G \square K_2$.

**Proof.** The forward direction is immediate. Suppose that Hadwiger’s Conjecture holds for $G \square K_2$; that is, $\chi(G \square K_2) \leq \eta(G \square K_2)$. Let $H$ be a graph with at least one edge and $\chi(G) \geq \chi(H)$. Then $\chi(G \square H) = \chi(G) = \chi(G \square K_2)$ by Sabidussi’s Theorem \[43\]. Since $K_2$ is a subgraph of $H$, $\eta(G \square H) \geq \eta(G \square K_2)$. In summary,

$$\chi(G \square H) = \chi(G) = \chi(G \square K_2) \leq \eta(G \square K_2) \leq \eta(G \square H).$$

Hence Hadwiger’s Conjecture holds for $G \square H$. \[\square\]
Theorem [12.5] motivates studying $\eta(G \square K_2)$ in more detail. By (4), $\eta(G \square K_2)$ is tied to $\text{tw}(G)$, the treewidth of $G$. By a minimum-degree-greedy algorithm, $\chi(G) \leq \text{tw}(G) + 1$. Thus it is tempting to conjecture that the lower bound on $\eta(G \square K_2)$ from Lemma [10.4] can be strengthened to

$$\eta(G \square K_2) \geq \text{tw}(G) + 1.$$  

This would imply that for all graphs $G$ and $H$ both with at least one edge and $\chi(G) \geq \chi(H)$,

$$\chi(G \square H) = \chi(G) \leq \text{tw}(G) + 1 \leq \eta(G \square K_2) \leq \eta(G \square H);$$

that is, Hadwiger’s Conjecture holds for every non-trivial product. However, (7) is false.

Kloks and Bodlaender [31] proved that a random cubic graph on $n$ vertices has $\text{tw}(G) \geq \Omega(n)$ but $\eta(G \square K_2) \leq O(\sqrt{n})$ by Lemma [2.2].

We finish with some comments about Hadwiger’s Conjecture for $d$-dimensional products. In what follows $G_1, \ldots, G_d$ are graphs, each with at least one edge, such that $\chi(G_1) \geq \cdots \geq \chi(G_d)$. Thus $\chi(G_1 \square G_2 \square \cdots \square G_d) = \chi(G_1)$ by Sabidussi’s Theorem [43]. Observe that Theorem [12.5] generalises as follows: Hadwiger’s Conjecture holds for all $G_1 \square G_2 \square \cdots \square G_d$ if and only if it holds for $G_1 \square Q_{d-1}$. (Recall that $Q_d$ is the $d$-dimensional hypercube.) Finally we show that if Hadwiger’s Conjecture holds for all graphs, then a significantly stronger result holds for $d$-dimensional products. By (3) and Theorem [7.1]

$$\eta(G_1 \square G_2 \square \cdots \square G_d) \geq \eta(K_{\eta(G_1)} \square Q_{d-1})$$

$$\geq \eta(K_{\eta(G_1)} \square K_{2^{d/2}})$$

$$\geq \eta(G_1) \cdot 2^{(d-2)/4} - O(\eta(G_1) + 2^{d/4})$$

$$\geq \chi(G_1) \cdot 2^{(d-2)/4} - O(\eta(G_1) + 2^{d/4})$$

$$= \chi(G_1 \square G_2 \square \cdots \square G_d) \cdot 2^{(d-2)/4} - O(\eta(G_1) + 2^{d/4}).$$

This shows that if Hadwiger’s Conjecture holds for all graphs, then the multiplicative factor of 1 in Hadwiger’s Conjecture can be improved to an exponential in $d$ for $d$-dimensional products.

**Recent Advances**

Chandran et al. [4] recently improved both the lower and upper bound in Theorem [7.1] to conclude that $\eta(K_n \square K_m) = (1 - o(1)) n \sqrt{m}$ for all $n \geq m$.

Lemma [10.4] can be restated as: if $\text{tw}(G) \geq 2^{4\ell^4}$ then $\eta(G \square K_2) \geq \ell$. This exponential bound was recently improved by Reed and Wood [39] to the following polynomial bound: for some constant $c$, if $\text{tw}(G) \geq c\ell^4 \sqrt{\log \ell}$ then $\eta(G \square K_2) \geq \ell$. Subsequently, the bounds in [4], in Theorem [12.3] and in the proof of Lemma [11.6] can be improved.
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