Multi-scale Stochastic Organization-Oriented Coarse-Graining
Exemplified on the Human Mitotic Checkpoint

– Supplementary Material –

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1 Reaction Network Models (Model 1 - Model 3)

Model 1 as rea-file format.
For the modified “short timescale” version (Model1b) we remove two reactions: “KinU -> KinA” and “O-Mad2 ->”. The modified version represents a short timescale at which no kinetochore gets attached and O-Mad2 does not decay.

# Number of Components
14
# Components
KinA
KinU
Apc
Mcc
Apc.Mcc
BubR1.Bub3
Cdc20
Apc.Cdc20
Cdc20.Mad2
C-Mad2
O-Mad2
Cdc20.BubR1.Bub3
Apc.Cdc20.C-Mad2
Apc.Cdc20.BubR1.Bub3
# Number of Reactions
21
# Reactions
1.0 KinU -> 1.0 KinA
1.0 Apc 1.0 Mcc -> 1.0 Apc.Mcc
1.0 Apc.Mcc -> 1.0 Apc 1.0 Mcc
1.0 Apc 1.0 Cdc20 -> 1.0 Apc.Cdc20
1.0 KinU 1.0 O-Mad2 -> 1.0 KinU 1.0 C-Mad2
1.0 Cdc20.Mad2 1.0 BubR1.Bub3 -> 1.0 Mcc
1.0 Mcc -> 1.0 Cdc20.Mad2 1.0 BubR1.Bub3
1.0 Cdc20 1.0 C-Mad2 -> 1.0 Cdc20.Mad2
1.0 C-Mad2 -> 1.0 O-Mad2
1.0 Cdc20.Mad2 -> 1.0 Cdc20 1.0 O-Mad2
1.0 Apc.Cdc20 1.0 Mcc -> 1.0 Apc 1.0 Cdc20 1.0 Mcc
1.0 Cdc20 1.0 BubR1.Bub3 -> 1.0 Cdc20.BubR1.Bub3
1.0 Cdc20.BubR1.Bub3 -> 1.0 Cdc20 1.0 BubR1.Bub3
1.0 Apc.Cdc20 1.0 C-Mad2 -> 1.0 Apc.Cdc20.C-Mad2
1.0 Apc.Cdc20.C-Mad2 -> 1.0 Apc.Cdc20 1.0 C-Mad2
1.0 Apc.Cdc20.C-Mad2 1.0 BubR1.Bub3 -> 1.0 Apc.Mcc
1.0 Apc.Mcc -> 1.0 Apc.Cdc20.C-Mad2 1.0 BubR1.Bub3
1.0 Apc.Mcc 1.0 KinA -> 1.0 Apc.Cdc20.BubR1.Bub3 1.0 O-Mad2 1.0 KinA
1.0 Apc.Cdc20.BubR1.Bub3 -> 1.0 Apc.Cdc20 1.0 BubR1.Bub3
1.0 O-Mad2 ->
1.0 Mcc 1.0 KinA -> 1.0 Cdc20.BubR1.Bub3 1.0 O-Mad2 1.0 KinA
### Model 2 as rea-file format.
For the modified “short timescale” version (Model2b) we remove reaction: “KinU -> KinA”.

| # Number of Components | 7 |
|------------------------|---|
| # Components           | Kin_A, Kin_U, Activator, Promotor, Promotor_Activated, Inhibitor, Promotor_Inactive |
| # Number of Reactions  | 9 |
| # Reactions            | 1.0 Kin_U -> 1.0 Kin_A, 1.0 Inhibitor -> 1.0 Activator, 1.0 Activator 1.0 Kin_U -> 1.0 Inhibitor 1.0 Kin_U, 1.0 Promotor_Inactive 1.0 Kin_A -> 1.0 Promotor_Activated 1.0 Kin_A, 1.0 Promotor 1.0 Activator -> 1.0 Promotor_Activated, 1.0 Promotor_Activated 1.0 Inhibitor -> 1.0 Promotor 1.0 Activator 1.0 Inhibitor, 1.0 Promotor 1.0 Inhibitor -> 1.0 Promotor_Inactive, 1.0 Promotor_Inactive -> 1.0 Promotor 1.0 Inhibitor, 1.0 Promotor_Activated -> 1.0 Promotor 1.0 Activator |

### Model 3 as rea-file format.
For the modified “short timescale” version (Model3b) we remove reaction: “KinU -> KinA”.

| # Number of Components | 4 |
|------------------------|---|
| # Components           | Activator, Inhibitor, Kin_A, Kin_U |
| # Number of Reactions  | 3 |
| # Reactions            | 1.0 Kin_U -> 1.0 Kin_A, 1.0 Activator 1.0 Kin_U -> 1.0 Inhibitor 1.0 Kin_U, 1.0 Inhibitor -> 1.0 Activator |
2 Supplementary Figure S1: State Transition Graph of a CTMC Model 3b)

Figure 1. State transition graph of the CTMC model for Model3b generated by PRISM. State labels show index and population count, e.g., 24 (0,1,0,3) denotes that there are 0 KinU, 1 KinA, 0 A and 3 I in state 24. Arrows denote the transitions between states, numbers over the arrows denote the rate of the transition.
3 Method: Approximate Aggregation Method for Model Reduction

In this section we recall a method\textsuperscript{24} for finding approximate aggregations of an ODE system $\dot{x} = A(x)$ where $A: \mathbb{R}^n \rightarrow \mathbb{R}^n$ is a quadratic polynomial. This case is important as any set of interesting biochemical reaction will usually incorporate rules for producing one type of particles by combining two different types of reactants. This kind of dynamics translates to $A$ being a polynomial map of degree two. Oddly enough, reactions involving terms of order 3 and higher are rarely encountered as they can, at least conceptually, be simulated as successive reactions of order two.

Firstly, recall that finding aggregations of the system $\dot{x} = A(x)$ is equivalent to finding aggregations of the map $x \mapsto A(x)$\textsuperscript{43}. It can be checked that any quadratic polynomial $A: \mathbb{R}^n \rightarrow \mathbb{R}^n$ can be written as

$$A(x) = A(0) + DA(0)x + \frac{1}{2}(DA(x) - DA(0))x$$

where $DA$ denotes the differential of $A$. Note that the second order term in the equation above is indeed linear as the other power of $x$ is hidden within the differential. It is possible to write our function in this convenient form only because we assumed that $A(x)$ is a quadratic polynomial and this sort of argument does not straightforwardly generalize to polynomials of higher degrees. Since entries of $DA$ are affine maps we can further write

$$DA(x) = DA(0) + \sum_{i=1}^{n} x_i(DA(e_i) - DA(0))$$

where $e_i$ for $1 \leq i \leq n$ are unit vectors of the canonical basis in $\mathbb{R}^n$ and $x$ is a vector with entries $x = (x_1, \ldots, x_n)$. Thus,

$$A(x) = A(0) + DA(0)x + \frac{1}{2} \sum_{i=1}^{n} x_i(DA(e_i) - DA(0))x$$

and we see that vector $A(x)$ can be written as a linear combination of vectors $A(0), DA(0)x, DA(e_1)x, \ldots, DA(e_n)x$. We will now show that any aggregation which is simultaneously compatible with matrices $DA(0), DA(e_1), \ldots, DA(e_n)$ must also be compatible with $A$. To see this, let us denote by $\Xi$ the coarse graining matrix associated to a particular aggregation which is simultaneously compatible with matrices $DA(0), DA(e_1), \ldots, DA(e_n)$. This means that we can find matrices $DA(0), DA(e_1), \ldots, DA(e_n)$ for which the following equations hold:

$$\Xi DA(0) = \overline{DA(0)}\Xi, \quad \Xi DA(e_1) = \overline{DA(e_1)}\Xi, \quad \ldots, \quad \Xi DA(e_n) = \overline{DA(e_n)}\Xi.$$

Then applying $\Xi$ to $A(x)$ we get:

$$\Xi A(x) = \Xi A(0) + \Xi DA(0)x + \frac{1}{2} \sum_{i=1}^{n} x_i(\Xi DA(e_i) - \Xi DA(0))x$$

$$= \Xi A(0) + + \overline{DA(0)}\Xi x + \frac{1}{2} \sum_{i=1}^{n} x_i(\overline{DA(e_i)} - \overline{DA(0)})\Xi x$$

where the last expression is clearly a function of $\Xi x$ which further means that $\Xi$ is compatible with map $A(x)$.

The converse to this fact is also known to be true\textsuperscript{45}. Together, this allows us to reduce our problem of finding aggregations for $A$ to that of finding aggregations simultaneously compatible with each of the matrices $DA(0), DA(e_1), \ldots, DA(e_n)$.

In fact, rather than looking for aggregations compatible simultaneously with $DA(0), DA(e_1), \ldots, DA(e_n)$ which may not even exist, it seems more reasonable to look for a set of matrices $\overline{DA(0)}, \overline{DA(e_1)}, \ldots, \overline{DA(e_n)}$ which approximate them in some matrix norm while at the same time being simultaneously compatible with a (preferably large) set of aggregations $\mathcal{S}$. The first step in doing so would be to do this for just one $n$-by-$n$ matrix $M$.

Recall that an $m$-by-$n$ ($m \leq n$) 0-1 matrix $\Xi$ with exactly one entry equal to 1 in each column is called an aggregation matrix and these are in 1-1 correspondence with the partitions of the set $\{1, 2, \ldots, n\}$. For each such $\Xi$ let $V_\Xi$ be the set of all $n$-by-$n$ matrices coarse grained by $\Xi$. One can show that this set is a linear subspace of the set of all matrices of order $n$ which we denote by $M_n(\mathbb{R})$\textsuperscript{24}. Thus, after fixing some matrix norm, one can orthogonally project matrix $M$ onto $V_\Xi$ in order to obtain the best approximation $\overline{M}$ which is coarse grained by $\Xi$.

But this only ensures that $\overline{M}$ is coarse grained by $\Xi$, whereas one would like to get as many coarse-grainings as possible. Note however that if $\Xi_1$ and $\Xi_2$ are two aggregation matrices, then $V_{\Xi_1} \cap V_{\Xi_2}$ is again a linear subspace of $M_n(\mathbb{R})$ consisting precisely of matrices which are coarse grained by both $\Xi_1$ and $\Xi_2$. Projecting orthogonally onto this subspace will yield an approximation of $\overline{M}$ which has at least those two valid reductions.

This idea extends to an arbitrary number of aggregation matrices, and one is inclined to ask how large a subset of aggregation matrices $\{\Xi_1, \ldots, \Xi_r\} \subseteq \{\Xi_0\}$ $|\, \Xi_0$ is a partition of $\{1, 2, \ldots, n\}$ can be while ensuring that the distance from $A$ to its orthogonal projection onto $V_{\Xi_1} \cap \cdots \cap V_{\Xi_r}$ is kept within a given error threshold. Note that the number of partitions of
a set \( \{1, \ldots, n\} \) is already super-exponential in \( n \), and going through all subsets of those would yield a growth rate greater than doubly exponential, thus making any brute force approach ineffective.

Instead, in Algorithm 1 below we propose running through all aggregations and retaining in a set \( S \) only those for which induced projections produce an approximation matrix within the error threshold \( \varepsilon \). The final approximation is then obtained by projecting onto \( \bigcap_{p \in S} V_p \). The reasoning behind this is that a matrix \( M \) that is \( \varepsilon \)-close to a subspace \( V \cap W \) cannot be further than \( \varepsilon \) from either \( V \) or \( W \). This however is not a guarantee that the converse, which is needed here, is true.

**Algorithm 1** Finding a nearby matrix that can be coarse grained

**Input:** (\( n \)-by-\( n \) matrix) \( M \)  
(error threshold) \( \varepsilon > 0 \)

\[
S = \{\} 
\]

for partition \( p \) of \( \{1, 2, \ldots, n\} \) do

\[
M' = \text{projection of } M \text{ onto } V_{\Xi_p} 
\]

if \( \|M - M'\| < \varepsilon \) then

add \( p \) to \( S \)

end if

end for

\[
V = \bigcap_{p \in S} V_p 
\]

\[
\tilde{M} = \text{projection of } M \text{ onto } V 
\]

**Output:** (the approximate matrix) \( \tilde{M} \)  
(the set of aggregations) \( S \)

It is not hard to modify Algorithm 1 so that it works with a set of matrices. We give the modification below as Algorithm 2. Again, there is no guarantee that the resulting matrices are \( \varepsilon \) close to the initial ones, but the same argument as before justifies this approach.

**Algorithm 2** Finding approximate matrices that can be jointly coarse grained

**Input:** (\( n \)-by-\( n \) matrices) \( M_1, \ldots, M_k \)  
(error threshold) \( \varepsilon > 0 \)

\[
S = \{\} 
\]

for partition \( p \) of \( \{1, 2, \ldots, n\} \) do

\[
M'_i = \text{projection of } M_i \text{ onto } V_{\Xi_p} 
\]

if \( \|M_i - M'_i\| < \varepsilon \), for all \( 1 \leq i \leq k \) then

add \( p \) to \( S \)

end if

end for

\[
V = \bigcap_{p \in S} V_{\Xi_p} 
\]

\[
\tilde{M}_i = \text{projection of } M_i \text{ onto } V 
\]

**Output:** (the approximate matrices) \( \tilde{M}_1, \ldots, \tilde{M}_k \)  
(the set of aggregations) \( S \)

Going back to our original problem, assume that we used this algorithm to obtain matrices \( \tilde{D}\!A(0), \tilde{D}\!A(e_1), \ldots, \tilde{D}\!A(e_n) \) and the set of aggregations \( S \) that simultaneously coarse grains each of them. Then the approximate quadratic map \( \tilde{A}(x) \) is simply given by

\[
\tilde{A}(x) = A(0) + \tilde{D}\!A(0)x + \frac{1}{2} \sum_{i=1}^{n} x_i (\tilde{D}\!A(e_i) - \tilde{D}\!A(0))x 
\]

and the ODE system \( \dot{x} = \tilde{A}(x) \) can be aggregated using any of the coarse grainings from the set \( S \).

### 3.1 Relation of this method to other model reduction techniques

The problem that standard model reduction techniques are attempting to solve usually consists of finding the evolution of a quantity of interest whose dynamics is modeled by a high-dimensional ODE system. The goal is then to find a low-dimensional model in which the evolution of the said quantity will match its actual evolution as accurately as possible.\(^{46,47}\)

In our approach we focus our attention to the high dimensional model itself and from there we attempt to derive quantities whose evolution can be *exactly* computed using reduced models. The guiding idea here being that it should be possible to automatically derive certain “conservation laws” that govern the dynamics of our system. The goal of our
approach is learning a hierarchical structure of the system under consideration and seeing how its organizational structure fits together. The model reduction part then comes as a by-product of this process.

Consequently, the approximate aggregation method discussed above favors models with rich structure, models with abundance of coarse grainings. This is perhaps best seen on an extremely simple example. The following two-dimensional linear ODE model is a modification of an example by Rowe et al.\textsuperscript{44}

\[
\begin{pmatrix}
\dot{x} \\
\dot{y}
\end{pmatrix} =
\begin{pmatrix}
.01 & .99 \\
.98 & .02
\end{pmatrix}
\begin{pmatrix}
x \\
y
\end{pmatrix}
\]

It can be calculated that both singular values of this matrix are of magnitude close to 1. The classical model reduction techniques based on SVD decomposition would therefore fail to reduce this model since both singular values are non-negligible and cutting off the smaller one would lead to a vastly different model.

\[
\begin{pmatrix}
\dot{x} \\
\dot{y}
\end{pmatrix} =
\begin{pmatrix}
.472 & .659 \\
.342 & .477
\end{pmatrix}
\begin{pmatrix}
x \\
y
\end{pmatrix}
\]

Our method on the contrary produces a system close to the original

\[
\begin{pmatrix}
\dot{x} \\
\dot{y}
\end{pmatrix} =
\begin{pmatrix}
.015 & .985 \\
.985 & .015
\end{pmatrix}
\begin{pmatrix}
x \\
y
\end{pmatrix}
\]

but in which variables \(x\) and \(y\) can be aggregated together to give a reduced model

\[
\dot{z} = z
\]

where \(z = x + y\).

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