STBC-Aided Cooperative NOMA with Timing Offsets, Imperfect Successive Interference Cancellation, and Imperfect Channel State Information

Muhammad Waseem Akhtar, Student Member, IEEE, Syed Ali Hassan, Senior Member, IEEE, Sajid Saleem, Member, IEEE, Haejoon Jung, Member, IEEE

Abstract—The combination of non-orthogonal multiple access (NOMA) and cooperative communications can be a suitable solution for fifth generation (5G) and beyond 5G (B5G) wireless systems with massive connectivity, because it can provide higher spectral efficiency, lower energy consumption, and improved fairness compared to the non-cooperative NOMA. However, the receiver complexity in the conventional cooperative NOMA increases with increasing number of users owing to successive interference cancellation (SIC) at each user. Space time block code-aided cooperative NOMA (STBC-CNOMA) offers less numbers of SIC as compared to that of conventional cooperative NOMA. In this paper, we evaluate the performance of STBC-CNOMA under practical challenges such as imperfect SIC, imperfect timing synchronization between distributed cooperating users, and imperfect channel state information (CSI). We derive closed-form expressions of the received signals in the presence of such realistic impairments and then use them to evaluate outage probability. Further, we provide intuitive insights into the impact of each impairment on the outage performance through asymptotic analysis at high transmit signal-to-noise ratio. We also compare the complexity of STBC-CNOMA with existing cooperative NOMA protocols for a given number of users. In addition, through analysis and simulation, we observe that the impact of the imperfect SIC on the outage performance of STBC-CNOMA is more significant compared to the other two imperfections. Therefore, considering the smaller number of SIC in STBC-CNOMA compared to the other cooperative NOMA protocols, STBC-CNOMA is an effective solution to achieve high reliability for the same SIC imperfection condition.

Index Terms—STBC, NOMA, cooperative NOMA, SIC, timing offset.

I. INTRODUCTION

Non-orthogonal multiple access (NOMA) is considered to be one of the most promising techniques for fifth-generation (5G) and beyond 5G (B5G) wireless systems to meet the heterogeneous demands on low latency, high reliability, massive connectivity, improved fairness, and high throughput [1]. The key principle behind NOMA is to exploit non-orthogonal resource allocation among multiple users at the cost of increased receiver complexity, which is required for separating the non-orthogonal signals [2]. In contrast to orthogonal multiple access (OMA), multiple users in NOMA are assigned the same physical resource (e.g., frequency and time) but with different power, which significantly enhances spectral efficiency.

Motivated by such advantage of NOMA, various aspects of NOMA have been actively investigated, engaging industry, standardization bodies, and academia. Further, as noted in [3], NOMA can be flexibly combined with various existing and emerging wireless technologies. In particular, the combination of NOMA and cooperative communications can be a suitable solution for the Internet-of-Things (IoT) networks with massive connectivity, because it can provide higher spectral efficiency, lower energy consumption, and improved fairness compared to the non-cooperative NOMA [4].

In one of the pioneering studies on the NOMA schemes that incorporates the principles of cooperative communications, the authors in [5] propose cooperative NOMA, which is subsequently referred to as conventional cooperative NOMA (CCN). In this scheme, strong users with better channel conditions support weak users with worse channel conditions by serving as relays, which increases the reliability of the weak users through diversity gain. In [6], cooperative NOMA is combined with simultaneous wireless information and power transfer (SWIPT) to improve energy efficiency through energy harvesting. Further, in [7], full duplex relaying-based NOMA schemes are introduced to reduce the number of time slots required to relay weak users’ messages. Similarly, cooperation among users by means of full-duplex device-to-device (D2D) communication is discussed in [8], where the outage performance of weak users is enhanced with the assistance of the full-duplex relaying by strong users. In addition, the authors in [9] propose a two-stage relay selection scheme for cooperative NOMA, which also provides lower outage rates. NOMA techniques adopting cooperative relaying systems are also extensively studied. For example, in [10], the authors propose an algorithm called cooperative relaying system using NOMA (CRS-NOMA), in which a decode-and-forward (DF) relay is adopted. Also, assuming a single DF relay and two far users, the outage performances of different relaying schemes are investigated in [11]. Also, NOMA using an amplify-and-forward (AF) protocol is investigated in [12].
Despite their effectiveness, the aforementioned NOMA schemes combined with cooperative communications incur an excessive number of successive interference cancellation (SIC) executed at user terminals as compared to the non-cooperative NOMA [13]. When it comes to the IoT networks with limited capabilities (e.g., computational resources and power), users may suffer from prohibitively large energy consumption due to the excessive number of SIC. To overcome this issue, space time block code (STBC)-aided cooperative NOMA protocols are proposed, which benefit from diversity gain with reduced number of SIC. For instance, the conventional Alamouti (i.e., $2 \times 1$) STBC-based NOMA system is investigated in [14], which uses two antennas at the BS and a single antenna at each user. This scheme doubles the diversity order as compared to that of conventional NOMA. Furthermore, in [15], the authors propose an Alamouti STBC-based CRS-NOMA protocol for a network with source, relay, and destination, which are equipped with two transmit antennas, two transmit and one receive antennas, and one receive antenna, respectively. It shows higher sum capacity and lower outage probability compared to the conventional CRS-NOMA in [10]. Instead of using the co-located (or real) antenna array, the authors in [13], [16] propose a new cooperative NOMA with a distributed STBC (i.e., STBC-CNOMA) for the virtual antenna array created by a group of single antenna users, which can be readily applied to the IoT networks. In their proposed scheme, STBC-CNOMA, they employ $2 \times 2$ distributed STBC on the NOMA system, in which two strong users act as DF relays and transmit the messages of the weak users by a $2 \times 2$ STBC. They show that STBC-CNOMA can achieve higher throughput with smaller number of SIC compared to the CCN in [5].

However, some challenges need to be addressed in the STBC-CNOMA systems in practical scenarios. For instance, distributed nature of terminals and their mobility cause the timing offsets, which is especially severe in virtual antenna array-based approaches including distributed STBC [17], [18]. In addition, reliability performance of NOMA can be significantly degraded by imperfect SIC, as reported in [19], and imperfect CSI [20]–[22]. However, the existing studies on STBC-CNOMA including [13] and [16] assume the ideal case without considering such realistic impairments. For this reason, in this paper, to better evaluate STBC-CNOMA, we investigate the impacts of the timing offsets, imperfect SIC, and imperfect CSI on its performance. Our main contributions can be summarized as follows.

- To the best of our knowledge, it is the first comprehensive study on practical impairments in STBC-CNOMA. We present a theoretical framework including signal model of an arbitrary user under the timing mismatch, imperfect SIC, and channel estimation error.
- We derive the probability distributions of the signal-to-interference-plus-noise ratio (SINR) for different combinations of the three impairments, which can be used in the STBC-CNOMA system design and operation. Based on the derived distributions, we also provide the closed-form expressions of outage probabilities, which are not present in prior arts.
- We also provide asymptotic rate (or capacity) outage probability in the high transmit signal-to-noise ratio (SNR) regime, which offers intuitive insights into how each of the three impairment hurts the performance of STBC-CNOMA.
- For the fair comparison with other cooperative NOMA protocols, we quantify the total number of SIC, the total number of required time slots, and the total number of transmissions of STBC-CNOMA with four existing schemes as functions of the number of user terminals.
- Numerical and simulation results are presented with different degrees of the three impairments. Through the comparison between analysis and simulation, we validate our analysis on the SINR distribution, exact and asymptotic capacity outage probabilities. In addition, we compare the outage performance of STBC-CNOMA with CCN and non-cooperative NOMA.

The rest of the paper is organized as follows. In Section II, we introduce the system model. Three practical impairments (i.e., timing error, imperfect SIC and imperfect CSI) and corresponding signal models are presented in Section III. Section IV provides closed-form expressions of outage rates both in the absence and in the presence of the imperfections. Furthermore, we compare the complexity of STBC-CNOMA with existing cooperative NOMA schemes in Section V. In Section VI, we present numerical and simulation results, and conclusions are drawn in Section VII.

**Notation:** $E[\cdot]$ and $\text{Var}[\cdot]$ denote the statistical expectation and variance, respectively. In addition, $|\cdot|$ denotes the absolute value of a scalar quantity. Also, the definitions of the variables used in our analysis are provided in Table I.

II. SYSTEM MODEL

We consider an STBC-based downlink NOMA system as shown in Fig. 1. Base station (BS) transmits the superimposed signal to all users in its coverage area. We assume the channel between the BS and the users and that between any two users
Table I: Table of Notations

| Symbol | Description |
|--------|-------------|
| $h_k$  | Channel gain from BS to the $k^{th}$ user |
| $g_k$  | Channel gain between the BS and the $k^{th}$ user |
| $\gamma_{k,noma}$ | SINR at the $k^{th}$ user to detect its own signal in direct NOMA phase |
| $\gamma_{k,ccn}$ | SINR at the $k^{th}$ user for conventional cooperative NOMA case [5] |
| $\gamma_{k,tl}$ | SINR at the $k^{th}$ user with perfect timing synchronization, perfect SIC (pSIC), and perfect CSI (pCSI) |
| $\gamma_{k,ipSIC}$ | SINR at the $k^{th}$ user with perfect timing synchronization, imperfect SIC (ipSIC), and perfect CSI (pCSI) |
| $\gamma_{k,ipCSI}$ | SINR at the $k^{th}$ user with imperfect timing synchronization, perfect SIC (pSIC), and perfect CSI (pCSI) |
| $\gamma_{k,ipSIC}$ | SINR at the $k^{th}$ user with imperfect timing synchronization, imperfect SIC (ipSIC), and perfect CSI (pCSI) |
| $\gamma_{k,ipCSI}$ | SINR at the $k^{th}$ user with perfect timing synchronization, perfect SIC (pSIC), and imperfect CSI (ipCSI) |
| $\gamma_{th}$ | SINR threshold |
| $r$    | Rate threshold |
| $p_k$  | Power received at the $k^{th}$ user from BS |
| $p_s$  | Mean power received by the user during STBC cooperation phase |
| $\xi_{k,t}$ | Noise received at the $k^{th}$ user during time slot $t$ of the STBC cooperation phase |
| $r_{k,t}$ | Received signal at the $k^{th}$ user during time slot $t$ of the STBC cooperation phase |
| $\lambda_A$ | Fading parameter for exponentially distributed variable $A$ |
| $\lambda_B$ | Fading parameter for exponentially distributed variable $B$ |
| $\lambda_F$ | Fading parameter for Gamma distributed variable $F$ |
| $\lambda_Z$ | Fading parameter for Gamma distributed variable $Z$ |
| erfc(x) | Error function |
| $E_i(x)$ | Exponential integral of $x$ and $E_i(x) = \int_{\infty}^{x} \frac{e^{-t}}{t} dt$ |

Transmission from BS to the users takes place in two phases. In the first phase, called the direct NOMA phase, BS sends the superimposed signal to all users. The weakest user extracts its own signal by considering the signals for all the other users as noise. Other users employ SIC to cancel the interference from the weak users and treat the signals for other strong users as noise. In the second phase, referred to as cooperative NOMA phase, the first two strongest users, $U_1$ and $U_2$, make an STBC pair and transmit the messages of the next two users, $U_3$ and $U_4$, by a distributed $2 \times 2$ STBC transmission. This process of $2 \times 2$ STBC continues until the weakest user $U_K$ is reached.

A. Direct NOMA Phase

As shown in Fig. 2, the direct NOMA phase is accomplished in the first time slot, when the BS transmits the superimposed signal to all of the $K$ users. The $k^{th}$ user, such that $1 \leq k < K$, detects the message of the $i^{th}$ user, where $i > k$, then applies SIC to subtract it from the superimposed signal and finally detects its own message.

B. STBC-based Cooperative Transmission Phase

The second phase of the proposed transmission is the STBC-based cooperative transmission phase. Users are paired as per their channel conditions, i.e., the first two strongest users make the first user pair followed by $U_3$ and $U_4$ making the second user pair until $U_{K-1}$ and $U_K$ making the $M^{th}$ user pair, where $M = K/2$ and $K$ is even. For the case $K$ is odd, $U_{K-1}$ and $U_K$ construct the $M^{th}$ user pair, where $M = \frac{K-1}{2}$. In this phase, all of the users cooperate with each other by employing a distributed $2 \times 2$ STBC transmission. However, at the receiving STBC users’ pair, we use $2 \times 1$ STBC reception for the detection of the symbols [23]. Each receiving STBC user receives two symbols, one for itself and the other for its neighbor. Thus, a STBC user keeps the decoded symbol for itself, while the other symbol for its neighbor is discarded. In the first time slot, $t_0$, BS transmits the composite NOMA signal to all users in its coverage area. Since the first two strongest users have decoded the messages for all the other users, therefore, they can contribute in the STBC cooperation by transmitting the information for next two users in the next two time slots. Therefore, during the second and third time
slots, $t_1$ and $t_2$, $U_1$ and $U_2$ transmit to the users $U_3$ and $U_4$ using Alamouti code. Similarly, in the two following time slots, $U_3$ and $U_4$ transmit the STBC signal to $U_5$ and $U_6$ and this process continues till $U_K$ receives its message. Assuming The SINR at each user of the receiving STBC pair, with perfect timing synchronization and perfect SIC, is given by

$$\gamma_k = \frac{|h_k|^2 p_k}{\sum_{i=1}^{k-1} |h_i|^2 p_i + \sigma^2} + \frac{(|g_k| g_{k-1}|^2 + |g_k| g_{k-2}|^2)p_s}{\sigma^2}, \quad (1)$$

where $2 < k \leq K$, $p_k = \Phi_k P_{NOMA}$ is the power assigned to the $k^{th}$ user, $\Phi_k$ is the power coefficient for the $k^{th}$ user, $P_{NOMA}$ is the power assigned to the composite NOMA signal and $p_s$ is the fraction of power transmitted from the transmitting users’ pair in STBC cooperation. Also, $i \in \{0,1\}$ denotes the first and the second user of the $2 \times 2$ STBC receiving pair, respectively. In case of conventional cooperative NOMA [5], each strong user, $U_i$, for any $i < k$ will cooperate with the weak user, $U_k$, by means of decode and forward relay. Assuming maximum-ratio combining (MRC), as in [5], the SINR at each user in this case is given by

$$\gamma_{k,ccn} = \frac{|h_k|^2 p_k}{\sum_{i=1}^{k-1} |h_i|^2 p_i + \sigma^2} + \frac{1}{\sum_{i=1}^{k-1} |g_k| |g_{k-j}|^2 + \sigma^2} g_{k,j}, \quad (2)$$

where $\gamma_{k,ccn}$ is the SINR at the $k^{th}$ user for conventional cooperative NOMA and $g_{k,j}$ is the power transmitted from the $(k-j)^{th}$ user to the $k^{th}$ user in the cooperation phase.

III. THREE PRACTICAL IMPAIRMENTS

In this section, we analyze the STBC-CNOMA system in the presence of the timing mismatch (or synchronization error) in the STBC cooperation phase, imperfect SIC, and channel estimation error. We treat the three imperfections in the three subsections, separately. To better explain the impacts of the three practical impairments, we use a simple example with four users (i.e., $K = 4$) as shown in Fig. 2. Using this example, we analyze the STBC transmission from $U_1$ and $U_2$ to $U_3$ and $U_4$. Then, based on the same approach, we extend our analysis to a general scenario with more number of users (i.e., $K > 4$).

A. Synchronization Error

Fig. 2 shows an STBC-based downlink NOMA network for two user pairs. During the first time slot $t_1$, each user receives the composite NOMA signal $X_{noma}$ from the base station. Since $U_1$ and $U_2$ are located in close vicinity of the BS, they decode their own messages in addition to the messages of $U_3$ and $U_4$ and send these to $U_3$ and $U_4$ through STBC transmission. In other words, $U_1$ and $U_2$ send $x_3$ and $x_4$ to $U_3$ and $U_4$ during time slot $t_1$. During next time slot $t_2$, $U_1$ and $U_2$ send $-x_3^*$ and $x_4^*$, respectively, to $U_3$ and $U_4$. Thus, the STBC receiving user pair $U_3$ and $U_4$ can detect their respective messages.

Fig. 3 illustrates the timing diagram with different synchronization conditions. Fig. 3(a) depicts the STBC mechanism at the receiver with perfect timing synchronization. The symbols from both users $U_1$ and $U_2$ of STBC pair arrive at the receiver at same time instant, where $T$ is the symbol duration. On the other hand, Fig. 3(b) shows the STBC transmission with timing offset of $\tau = \varepsilon_2T$. In this case, the symbols from $U_1$ and $U_2$ does not arrive simultaneously and there is a substantial inter-symbol-interference (ISI) experienced by the STBC receiving user pair, which causes the decrease in the SINR. The STBC block used by $U_1$ and $U_2$ is given by

$$S = \begin{bmatrix} x_3 & -x_4^* \\ x_4 & x_3^* \end{bmatrix}. \quad (3)$$

Thus, if $U_4$ is not perfectly synchronized, as shown in Fig. 3(b), the receiver equations for $U_3$ or $U_4$ of STBC receiving pair at time slots $t_1$ and $t_2$ are given as

$$r_{k,1} = g_{3,1} x_3 + g_{3,2} e_1 x_4 + \xi_k, \quad (4)$$
$$r_{k,2} = -g_{3,1} x_3^* + g_{3,2} e_1 x_4^* + g_{3,2} e_2 x_4 + \xi_k, \quad (5)$$
$$r_{k,3} = g_{4,1} x_3 + g_{4,2} e_1 x_4 + \xi_k, \quad (6)$$
$$r_{k,4} = -g_{4,1} x_3^* + g_{4,2} e_1 x_4^* + g_{4,2} e_2 x_4 + \xi_k, \quad (7)$$

where $r_{k,t}$ and $\xi_{k,t}$ are the received signal and the additive noise observed at the $k^{th}$ user during time slot $t$, respectively. The received signals after combiner at User 3 and User 4 are given by

$$\tilde{v}_3 = g_{3,1} x_3 + g_{3,2} x_4 + \xi_3, \quad (8)$$
$$\tilde{v}_4 = g_{4,1} x_3^* + g_{4,2} x_4^* + \xi_4, \quad (9)$$

which can be expanded as

$$\tilde{v}_3 = (|g_{3,1}|^2 + e_1|g_{3,2}|^2)x_3 + (e_1 - 1)g_{3,1} g_{3,2} x_4$$
$$+ e_2|g_{3,2}|^2 x_4^* + g_{3,1} e_3 + g_{3,2} e_3^*, \quad (10)$$

$$\tilde{v}_4 = (|g_{4,1}|^2 + e_1|g_{4,2}|^2)x_4 + (1 - e_1)g_{4,1} g_{4,2} x_3$$

From User 1 of STBC user Pair
From User 2 of STBC user Pair
From User 1 of STBC user Pair
From User 2 of STBC user Pair

Figure 3: Timing diagram of the received signals.
As a result, assuming MRC of the received signals in both direct NOMA and STBC phases, the SINRs of $U_3$ and $U_4$ can be obtained as (12) and (13), respectively. We note that if timing synchronization is perfect (i.e., $\varepsilon_1 = 1$ and $\varepsilon_2 = 0$), the corresponding SINRs in (12) and (13) are reduced into (1). Furthermore, generalizing this four-user example to a larger number of users (e.g., $K = 6, 8, \ldots$), in the presence of the synchronization error, the mathematical expressions of the SINR of the $m^{th}$ and the $n^{th}$ users can be derived as (14) and (15), respectively, for any $m \in \{3, 5, 7, \ldots, K-1\}$ and $n \in \{4, 6, 8, \ldots, K\}$.

### B. Imperfect SIC

Without synchronization error but under the residual interference caused by imperfect SIC implementation, the SINRs at the $m^{th}$ and the $n^{th}$ users are given by

$$
\gamma_m = \frac{|h_m|^2 p_m}{\eta |g_n|^2 p_n + \sum_{i=1}^{m-1} |h_m|^2 p_i + \sigma^2 + \frac{(|g_{m,m-1}|^2 + |g_{m,m-2}|^2)p_s}{\sigma^2}},
$$

and

$$
\gamma_n = \frac{|h_n|^2 p_n}{\eta |g_n|^2 p_n + \sum_{i=1}^{n-1} |h_n|^2 p_i + \sigma^2 + \frac{(|g_{n,n-1}|^2 + |g_{n,n-3}|^2)p_s}{\sigma^2}},
$$

respectively. Therefore, in the presence of both the synchronization error and imperfect SIC, the SINRs of the $m^{th}$ and the $n^{th}$ users can be derived as (16) and (17), respectively, where $m \geq 3$ and $n \geq 4$ can be any odd and even numbers, respectively. In addition, it is to be noted that $\eta = 0$ and $\eta = 1$ represent the perfect and imperfect SIC employed at that user, respectively. It is noted that the SIC imperfection and timing error do not affect each other.

### C. Imperfect CSI

In this section, we consider the imperfect CSI (i.e., channel estimation error) along with imperfect timing synchronization. We consider $\hat{g}_{k,j} = g_{k,j} + \omega_j$, where $\hat{g}_{k,j}$ is the estimate of $g_{k,j}$ and $g_{k,j}$ is the channel between $k^{th}$ and $j^{th}$ user. $\omega_j$ is the channel estimation error and it is assumed to be complex Gaussian random variable (RV) with zero mean and variance of $\sigma^2_g$. The RV $\hat{g}_{k,j}$ are complex Gaussian with zero mean and variance $\sigma^2_\omega = \sigma^2_g + \sigma^2_g$. Also, the correlation coefficient between the estimated channel and the real channel is $\rho = \sigma^2_\omega / (\sigma^2_g + \sigma^2_\omega)$. We can write that $g_{k,j} = \rho \hat{g}_{k,j} + \tilde{g}_{k,j}$, where $\tilde{g}_{k,j}$ are independent complex Gaussian RVs with zero mean and variance $\sigma^2_\tilde{g} = \sigma^2_\omega / (\sigma^2_g + \sigma^2_\omega)$ [20, 21]. By putting the value $g_{k,j} = \rho \hat{g}_{k,j} + \tilde{g}_{k,j}$ into (8) and (9) [24], we get

$$
\tilde{\nu}_4 = \psi_4 (g_{3,1} + \rho g_{3,1}) (g_{3,1} + \rho g_{3,1})^* + \psi_{4,1} (g_{3,2} + \rho g_{3,2})^* x_3,
$$

and

$$
\tilde{\nu}_4 = [(g_{4,2} + \rho g_{4,2}) (g_{4,2} + \rho g_{4,2})^* + (g_{4,1} + \rho g_{4,1}) (g_{4,1} + \rho g_{4,1})^*] x_4

\xi_{4,1} (g_{4,2} + \rho g_{4,2})^* - \xi_{4,2} (g_{4,1} + \rho g_{4,1})^*,
$$

respectively. Assuming the MRC of the received signals at each user for direct NOMA and STBC cooperation phase, and solving (21)-(22) for the SINRs, we obtain (18) and (19), respectively.

### IV. Outage Probability Analysis

In this section, we analyze the outage performance under the three practical impairments: timing error, imperfect SIC, and channel estimation error. Because the last user (i.e., User $K$) has the weakest channel gain and also suffers from the impairments, it has the worst outage probability compared to the other users, as shown in (13), (16). For this reason, we focus on the outage performance of User $K$ (e.g., User 4 in the four-user example in the previous section), which will set a benchmark for the other users with stronger channel gains.

An outage event occurs when a user cannot achieve the reliable SINR to detect the signal. Following (13) and (16), the outage probability of any user $k$ (for $k \leq K$) is defined as

$$
P_{out} = P(\gamma_k < \gamma_{th}) = \int_{0}^{\gamma_{th}} f_{\Gamma}(\gamma_k) d\gamma_k, \quad \gamma_{th} < 2^{Y} - 1,
$$

where $Y = \log_2(1 + \gamma_{th})$ is the rate threshold. In order to find the outage probability using (24), we need the PDF of SINR for different cases. Therefore, as the SINR expressions in (14)-(19) contain random variables, we consider the following mathematical manipulation and define some composite random variables for finding the PDF of SINRs for different cases.

We redefine the variables used in (14)-(19) as $A = |h_k|^2 p_k$, $B = \sum_{i=1}^{K} |h_k|^2 p_i$, $C = |g_{k,k-2}|^2 p_s$, $D = |h_{k,k-3}|^2 p_s$, $F = |g_{n,n-2}|^2 p_s$, $C^\chi = |A^\chi|^2 p_s$, $D^\chi = |B^\chi|^2 p_s$, $|h_k|^2$, $|g_{n,n-2}|^2$, $|g_{k,k-2}|^2$, $|g_{k,k-3}|^2$, $|A^\chi|^2$, $|B^\chi|^2$ follow the exponential distribution with parameters $\lambda_k$, $\lambda_{g_{n,n-2}}$, $\lambda_{g_{k,k-2}}$, $\lambda_{g_{k,k-3}}$ and $\lambda_{\chi}$, respectively. The variables $A$, $C$, $D$, $F$, $C^\chi$ and $D^\chi$ also follow the exponential distributions with parameters $\lambda_{\chi}$, $\lambda_{g_{k,k-2}}$, $\lambda_{g_{k,k-3}}$, $\lambda_{\chi}$, $\lambda_{\chi}$ and $\lambda_{\chi}$, respectively. We assume that $\lambda_{g_{k,k-2}} = \lambda_{g_{k,k-3}} = \lambda_{\chi}$ and $\lambda_{\chi} = \lambda_{\chi} = \lambda_{\chi}$, where disparate path losses with large-scale fading are compensated by appropriate power control at the relaying users, as in (25).

The variable $B$ follows the hypo-exponential distribution with
parameters $\lambda_i$, where $i \in \{1, 2, 3, \ldots, I\}$ is a set of interfering users and $I = k - 1$. We denote $\lambda_k = \frac{1}{p_\rho \gamma}, \lambda_i = \frac{1}{p_{i\rho} \gamma}$, $\lambda_y = \frac{1}{p_{i\rho} \gamma}$, and $\lambda_x = \frac{1}{p_{i\rho} \gamma}$, where $p_\eta$ and $p_\chi$ is the power of interfering signal (IS) due to imperfect SIC and power of IS due to imperfect CSI, respectively.

In the following five lemmas and five propositions, we treat different combinations of the three impairments. In the lemmas, we derive the exact outage probabilities based on $\overline{P}_{\text{out}}$ in (25). The probability distributions of the SINRs in each case can be found in its proof. Further, in the propositions, which correspond to each of the five lemmas, we provide the asymptotic outage probabilities in the high transmit SNR regime using $\overline{P}_{\text{out}}$ in (25), which provide intuitive insights into the impacts of the impairments. For mathematical notations used to derive the lemmas, please refer to Table III.

First, we consider the outage probability in the absence of any impairments, which can serve as a baseline to quantify the impact of the imperfections, in the following lemma.

**Lemma 1.** The outage probability for the perfect timing, perfect SIC, and perfect CSI is given as

$$P_{\text{out}} = \frac{1}{\lambda^2 \gamma^2} \sum_{i=1}^{I} \left[ \frac{\gamma_i}{\lambda_i} + \frac{\lambda_i}{\gamma_i} \right] e^{-\lambda_i - \lambda_i \gamma_{\text{th}}} e^{\lambda_i \gamma_{\text{th}} \left( \frac{\lambda_i}{\lambda_y \lambda_h} \right)}$$

where $\gamma_{\text{th}}$ is the SINR threshold and $\delta = \frac{1}{\lambda_k} \prod_{i=1}^{I} \frac{\lambda_i}{\lambda_k} \left( 1 - \frac{\lambda_i}{\lambda_k} \right)$.

**Proof:** See Appendix A.

Since the exact outage expression derived in Lemma 1 is complicated and does not yield to easy interpretation, we consider the asymptotic behavior of the rate outage given in (25), when the transmit SNR, which is denoted by $\gamma$, is high enough in the following proposition.

**Proposition 1.** As $\text{SNR} \to \infty$, the rate outage probability for the perfect timing synchronization, perfect SIC, and perfect CSI becomes

$$\lim_{\text{SNR} \to \infty} \overline{P}_{\text{out}} \sim \frac{\lambda_0 g(\text{SNR})}{\Phi_k - \sum_{i=1}^{I} \Phi_i (2\gamma - 1)} + 2\lambda_0^2 g(\text{SNR})^2,$$

where $g(\text{SNR}) = \frac{\gamma - 1}{\text{SNR}}$. 

(27)
Table II: Mathematical Notations

| Mathematical Notations used in Corollaries and Appendices |
|----------------------------------------------------------|
| \( \psi_i = \prod_{j=1}^{I} \lambda_j, \psi_i^n = \lambda_n \rho_n \prod_{j=1}^{I} \lambda_j, \psi_2 = \sum_{i=1}^{I} \prod_{j=ij} \lambda_j, \) |
| \( \psi_2^n = \lambda_n + \psi_2, \psi_3 = \sum_{i=1}^{I} \log[\lambda_n] \prod_{j=1,j \neq i}^{I} (\lambda_j - \lambda_k), \psi_4 = \lambda_n \psi_4, \) |
| \( \psi_4^n = \lambda_n \sum_{i=1}^{I} \log[\lambda_n] \prod_{j=1,j \neq i}^{I} (\lambda_j - \lambda_k) \prod_{j=1}^{I} (\lambda_j - \lambda_k), \) |
| \( \psi_4 = \lambda_n \sum_{i=1}^{I} \log[\lambda_n] \prod_{j=1,j \neq i}^{I} (\lambda_j - \lambda_k), \psi_6 = \sum_{i=1}^{I} \log[\lambda_n] \prod_{j=1}^{I} (\lambda_j - \lambda_k), \) |
| \( \psi_7 = \prod_{j=1,k=1}^{I} (\lambda_j - \lambda_k)^2, \psi_8 = \sum_{i=1}^{I} \lambda_i \log[\lambda_i] \prod_{j=1,j \neq i}^{I} (\lambda_j - \lambda_k), \) |
| \( \psi_9 = \lambda_n^2 \psi_9, \psi_{10} = \sum_{i=1}^{I} e^{\lambda_i} \prod_{j=1,j \neq i}^{I} (\lambda_j - \lambda_k) E_i(\lambda_k), \) |
| \( \psi_{10} = \sum_{i=1}^{I} \sum_{j=1}^{I} \sum_{s=1}^{I} \sum_{t=1}^{I} \lambda_s \lambda_t \lambda_i \lambda_j, \psi_i = \prod_{j=1,k=1}^{I} (\lambda_j - \lambda_k)^2 \) |

Proof: See Appendix B.

Based on this ideal case, we will investigate how each impairment impacts the outage probability.

Lemma 2. The outage probability for the perfect timing synchronization, imperfect SIC, and perfect CSI is given as

\[
P^{\eta}_{\text{out}} = \sum_{i=1}^{I} \left[ \frac{1}{\lambda_n^2 \lambda_h^2 (\lambda_i - \lambda_n)} \right] \Theta \left\{ \Theta \left( \frac{\lambda_n e^{-\lambda_n + \lambda_h + \gamma_{\text{th}} \lambda_n}}{\lambda_n^2 \lambda_h} \right) - \Theta \left( \frac{\lambda_n e^{-\lambda_n + \lambda_h + \gamma_{\text{th}} \lambda_n}}{\lambda_n} \right) \right\} \Theta \left( \frac{\lambda_n e^{-\lambda_n + \lambda_h + \gamma_{\text{th}} \lambda_n}}{\lambda_n} \right)
\]

where the values of \( \theta \) and \( \phi \) are given in (65).

Proof: See Appendix G.

We note that Lemmas 1, 2, and 3 are special cases of this lemma. The following asymptotic analysis provides the corresponding limiting rate outage for \( \text{SNR} \to \infty \).

Proposition 4. As \( \text{SNR} \to \infty \), the rate outage probability for the imperfect timing synchronization, imperfect SIC, and perfect CSI becomes

\[
\lim_{\text{SNR} \to \infty} \hat{P}^{\eta}_{\text{out}} \sim \hat{g}(\text{SNR}) + g^\phi(\text{SNR}).
\]
Proof : See Appendix H.

In fact, this result in \( (33) \) is in line with the composite degradations found in Propositions 2 and 3. In the following lemma and proposition, we will investigate the impact of the imperfect CSI.

**Lemma 5.** The outage probability for the perfect timing synchronization, perfect SIC, and imperfect CSI is given as

\[
P_{\text{out}}^{x} = \frac{1}{\lambda_{x}^{2} \lambda_{h}} \sum_{i=1}^{l} \left[ \Phi_{h} \left( \frac{\gamma_{x}}{\lambda_{h}^{2}} \right) + 2 \lambda_{x}^{2} \left( \frac{\gamma_{x}}{\lambda_{h}^{2}} \right)^{2} \right] + 2 \lambda_{x} \sum_{i=1}^{l} \Phi_{h} \left( \frac{\gamma_{x}}{\lambda_{h}^{2}} \right)
\]

Proof : See Appendix I.

**Proposition 5.** As \( \text{SNR} \to \infty \), the outage probability for the perfect timing synchronization, perfect SIC, and imperfect CSI becomes

\[
\lim_{\text{SNR} \to \infty} \tilde{P}_{\text{out}}^{x} \sim \frac{\lambda_{h} g(\text{SNR})}{\Phi_{k} - \sum_{i=1}^{K-1} \Phi_{i} (2^{T} - 1)} + 2 \lambda_{x}^{2} [g(\text{SNR})]^{2}, \quad (35)
\]

where \( \lambda_{x} = 1/\sigma_{x}^{2} = 1/(\sigma_{g}^{2} + \rho^{2} \sigma_{g}^{2}) \).

Proof : See Appendix J.

Since the outage rates are obtained in closed-form expressions in the five lemmas, it is possible to estimate how the outage performance of the STBC-CNOMA scheme changes as various system parameters change. In addition, the asymptotic analysis in the propositions provide the insights into how the rate outage is degraded by the three imperfections. Our analysis in the lemmas and propositions will be validated by comparing with simulation results in Section VI.

### V. Complexity Analysis

In this section, we will compare the complexity of STBC-aided cooperative NOMA (STBC-CNOMA) with other cooperative NOMA techniques in terms of the number of SIC performed, which is widely used to evaluate the computational complexity of NOMA as in \([6, 7, 26-28]\). We consider five different NOMA schemes: conventional cooperative NOMA (CCN) \([5]\), cooperative relay systems using NOMA (CRS-NOMA) \([10]\), CRS-NOMA novel design (CRS-NOMA-ND) \([29]\), cooperative relay selection by using STBC (CRS-STBC-NOMA) \([15]\), and lastly STBC-CNOMA. We note that this is the first extensive comparison of the five schemes, which quantifies the total number of SIC performed for a given number of users \( K \).

The total number of SIC performed by CCN is given by

\[
SIC_{\text{ccn}} = \sum_{j=0}^{K-2} \left[ \sum_{i=1}^{K-1-j} (K - (i + j)) \right]. \quad (36)
\]

Also, the total number of SIC performed by CRS-NOMA, which uses a half-duplex relay between the source and each user, is given by

\[
SIC_{\text{crs-noma}} = \sum_{j=1}^{K} j. \quad (37)
\]

Also, the total number of SIC performed by CRS-STBC-NOMA, which has a two-phase communication from the source to each user by means of 2 \( \times 1 \) STBC, is given by

\[
SIC_{\text{crs-stbc-noma}} = \sum_{j=1}^{K} 4j. \quad (38)
\]

In this scheme, the source is equipped with two transmit antennas, and each relay is equipped with one receive antenna and two transmit antennas, whereas each user is equipped with one receive antenna.

On the other hand, the total number of SIC performed by CRS-NOMA-ND is given by

\[
SIC_{\text{crs-noma-nd}} = \sum_{j=1}^{K} 2j. \quad (39)
\]

In this scheme, two sources transmit two symbols to two users by means of superposition coding, and each user decodes its symbol by MRC and SIC. Whereas, the total number of SIC performed by STBC-CNOMA is given by

\[
SIC_{\text{stbc-cnoma}} = \sum_{i=1}^{K-1} (K - i). \quad (40)
\]

It is noted that \( (36) \) and \( (40) \) show that for a larger number of users, complexity of CCN increases due to higher number of SIC. Whereas, the number of SIC is not increased in cooperation phase of STBC-CNOMA. Therefore, it achieves the diversity gain of STBC codes and maintains the number of SIC same as that of conventional direct NOMA scheme.

In Table III, we compare the number of time slots required for the complete transmission in various cooperative schemes for even \( K \geq 4 \). As shown in the table, STBC-CNOMA requires \( K - 1 \) time slots for the complete transmission. Therefore, saving one time slot as compared to that of in CCN and CRS-NOMA. This time slot can be utilized in other types of signaling, which makes STBC-CNOMA more efficient scheme. In addition, in the last column of Table III we compare the number of transmissions, which indicates the communication overhead. As shown in the table, CRS-STBC-NOMA has the largest number of transmissions.
function of means that the outage probability of the user is an increasing CCN, and STBC-CNOMA for τ of the user decreases, as the timing offset distribution approaches zero. In other words, the average SINR as expected, as the timing offset increases, the mean of the SINR PDFs derived in Section VI have been validated. Also, as an essential component for further performance analysis, the results for various timing offsets τ require more number of transmissions compared to the other three scheme for any K ≥ 4.

VI. Simulation Results

In this section, we present numerical and simulation results and validate our analysis in the previous sections. We consider a downlink NOMA system with one BS and K users. As in [5], [10], [13], we assume that the channels between the BS and each user, and inter-users are flat fading Rayleigh channels. The noise power spectral density is considered as −174 dBm/Hz. The rate threshold Υ is set to be 2 bits per channel use (BPCU). Each user is considered as stationary. The parameters ζ₀, ζ₁, ζ₂, ζ₃, ζ₄ and ζ₅ are considered to be unity. Symbol duration T is also considered as unity for the simplicity. For the STBC-CNOMA scheme, the transmit power at the BS (i.e., P_{NOMA}) is considered as 45 dBm, whereas the power transmitted during STBC cooperation phase (i.e., pₛ) is considered as half of the power transmitted from the BS. Power coefficients for simulations with K = 4 are Φ₁ = 0.1, Φ₂ = 0.2, Φ₃ = 0.3, and Φ₄ = 0.4 for k ∈ {1, 2, 3, 4}, respectively. It is to be noted that we assume the same total power budget for all of the schemes in the simulation for the fair comparison.

Fig. 4 provides the comparison of simulation and analytical results for the SINR PDFs for different timing offsets, assuming K = 4 with the perfect SIC and CSI. It can be noticed that the simulation results closely match the analytical results for various timing offsets τ ∈ {0, 0.3, 0.6, 0.9}. Thus, as an essential component for further performance analysis, the SINR PDFs derived in Section VI have been validated. Also, as expected, as the timing offset increases, the mean of the distribution approaches zero. In other words, the average SINR of the user decreases, as the timing offset τ increases, which means that the outage probability of the user is an increasing function of τ.

In Fig. 5, we compare the outage probabilities of NOMA, CCN, and STBC-CNOMA for τ = {0, 0.5, 0.8} and K = 4.

In the figure, the horizontal axis represents the SINR threshold γ_{th}. The simulation and analytical results show great correlation with each other, which validates Lemmas 1 and 3. Further, in the figure, we observe that the outage probability of NOMA is the highest for a given γ_{th}, because it is a non-cooperative scheme that does not provide diversity gain. It is also noted that STBC-CNOMA outperforms the CCN for the low SINR thresholds. Also, the performance of the STBC-CNOMA with τ = 0.5 is similar to that of CCN for SINR threshold of up to 2 dB. Fig. 5 also demonstrates that the outage probability of the STBC-CNOMA approaches to that of NOMA with the timing offset, τ approaching to 1. Therefore, for the low SINR threshold and τ < 0.5, STBC-CNOMA is still an attractive scheme.

Fig. 6 shows both analytical and simulation results of the
outage performance of the STBC-CNOMA with the total number of users \( K \in \{4, 8, 16\} \). In the figure, we observe the outage performance degradation with the increasing number of users. Further, it is noted that there is not much difference in the outage performance as \( K \) exceeds 8. In other words, in the absence of any imperfection, the outage performance of STBC-CNOMA is not much degraded with the increasing number of users. Also, the analytical and simulation results closely match, which validates Lemma 2 in Section IV.

Similarly, Fig. 7 shows the analytical and simulation results for the performance of STBC-CNOMA with different timing offsets for the perfect CSI and imperfect SIC of \(-5\) dBs, when \( K = 4 \). Fig. 8 depicts the outage probability with different levels of the SIC imperfection. In this case, Fig. 7 and Fig. 8 show the great correlation between the simulation and analytical results based on Lemma 4 in Section IV. Also, as shown in the figures, we observe that the outage probability increases sharply, as \( \tau \) and the SIC imperfection increase.

Fig. 9 presents a comparative analysis of the rate outage performance as function of the transmit SNR, \( \text{SNR} \), of User 4 for STBC-CNOMA with the perfect SIC and different values of timing offsets. We assume that the rate threshold is 2 bits per channel use (BPCU) and there are 4 users (i.e., \( K = 4 \)) in the system. Also, the dotted lines correspond to the asymptotic rate outage probabilities derived in Section IV. The results show that the rate outage degrades with the increase in the timing offset \( \tau \), which is in line with the previous simulation.

Figure 7: Outage probability performance with \( K = 4 \) for the imperfect SIC = -5 dBs, perfect CSI, and variable timing offsets.

Figure 8: Outage probability of User 4 in STBC-CNOMA with the imperfect SIC, perfect timing synchronization, and perfect CSI.

Figure 9: Rate outage probability of User 4 in STBC-CNOMA at \( \tau = \{0, 0.6, 0.7, 0.8, 0.9\} \), perfect CSI, and perfect SIC.

Figure 10: Rate outage probability of User 4 in STBC-CNOMA with the perfect timing synchronization, perfect CSI, and imperfect SIC.
Figure 11: Rate outage probability of User 4 in STBC-CNOMA with the timing offset of 0.5, perfect CSI, and imperfect SIC.

Figure 12: Rate outage probability of User 4 in STBC-CNOMA with the timing offset of 0.7, perfect CSI, and imperfect SIC.

Figure 13: Rate outage probability of User 4 in STBC-CNOMA with the perfect timing synchronization, perfect SIC, and imperfect CSI.

Figure 14: A comparison of number of SIC performed in different flavors of cooperative NOMA.
is desirable in hardware-limited networks.

The proposed scheme, STBC-CNOMA, outperforms all of the other schemes in terms of numbers of SIC performed, which is the case in cooperative NOMA protocols such as CCN, CRS-NOMA, CRS-STBC-NOMA and CRS-NOMA-ND in terms of the total number of SIC. Through both analysis and simulation, we show that STBC-CNOMA can be an attractive solution for systems with higher number of users or devices and having low power constraints. The simulation results also have shown that for a small number of users (i.e., \( K \leq 4 \)) STBC-CNOMA without any imperfection outperforms CCN, until the SINR threshold exceeds a certain value. Even with moderate timing offset \( \tau < 0.5 \), we have observed that the outage performance degradation of STBC-CNOMA relative to CCN is not significant. On the other hand, the impact of the imperfect SIC on the outage performance of STBC-CNOMA is more significant compared to those of the other cooperative NOMA protocols, we can conclude that STBC-CNOMA is an effective solution to achieve high reliability for the same SIC imperfection condition.

APPENDIX A: PROOF OF LEMMA 1

Suppose that \( A \sim \text{Exp}(\lambda_a) \) and \( B \) is the sum of exponential RVs resulting in a hypo-exponential RV, i.e., \( B \sim \text{hypoexp}(\lambda_4) \) with \( \lambda_4 \) is a vector given as \( \lambda_4 = \{\lambda_1, \lambda_2, \lambda_3, \ldots, \lambda_I\} \), as described in Section IV-A, then the PDF of \( A \) is given by

\[
f_A(a) = \lambda_h e^{-\lambda_h a}, \quad a \geq 0.
\]

In case of only one interfering user, the PDF of \( B \) is given as

\[
f_B(b) = \lambda_1 e^{-\lambda_1 b}, \quad b \geq 0.
\]

For more than one interferers, the PDF of \( B \) is given as

\[
f_B(b) = \left( \prod_{i=1}^I \lambda_i \right) \left( \sum_{i=1}^I \frac{e^{-\lambda_i b}}{\prod_{j=1, i \neq j}^I (\lambda_j - \lambda_i)} \right),
\]

where \( b > 0 \) and \( I \) is the total number of interferers. The PDF of \( Q_1 = \frac{A}{B} \) can be expressed as

\[
f_{Q_1}(q) = \int_0^\infty b f_B(b) f_{A}(bq) db.
\]

By substituting (41) and (42) into (44), we get the form in (45). Then, as in [30], the PDF of \( Q_1 = \frac{A}{B} \) is given by

\[
f_{Q_1}(q) = \frac{\lambda_h \psi_1}{\prod_{j=1}^I (\lambda_h q + \lambda_j)^2} \left[ \psi_2 + q \lambda_h \left( \sum_{i=1}^I (i+1) q^{i-1} \lambda_h^{i-1} \psi^{i-1} \right) \right],
\]

where \( q > 0 \). Hence, using this PDF, the mean and variance of \( Q_1 \) can be obtained as \( \mathbb{E}[Q_1] = \frac{\lambda_h \psi_1}{\psi_2} \) and \( \text{Var}[Q_1] = \frac{\lambda_h \psi_1 (\lambda_h \psi_1 \psi_2 + 2 \psi \psi_2)}{\psi_2^2} \), respectively. Similarly, for the perfect timing synchronization, let \( Z = C + D \), which is the sum of two exponential RVs. Thus, \( Z \sim \text{Gamma}(\lambda_2) \), which corresponds to the following PDF of \( Z \)

\[
f_Z(z) = \frac{z}{\lambda_2^2} e^{-z \lambda_2}, \quad z > 0.
\]

Without any imperfection, we can write (1) as \( L = Q_1 + Z \), where the PDF of \( Q_1 \) is given as (45). By convolving the

VII. CONCLUSIONS

In this paper, we have provided the theoretical framework to incorporate three realistic impairments, which are timing error, SIC imperfection, and CSI impairment, with outage performance. We have derived the closed-form expressions of the outage probabilities for the different combinations of the three impairments. Further, the complexity of STBC-CNOMA has been compared with existing cooperative NOMA protocols such as CCN, CRS-NOMA, CRS-STBC-NOMA and CRS-NOMA-ND in terms of the total number of SIC. Through
PDFs of \( Q_1 \) and \( Z \), we obtain the PDF of \( L = Q_1 + Z \), which can be found in [30] as
\[
f_L(l) = \frac{1}{\lambda_g^3} \sum_{i=1}^{I} \left[ 0e^{-\frac{\lambda_g g h}{\lambda_g + \lambda_i}} \left( \lambda_i (\lambda_g \lambda_h + \lambda_i + \lambda_h l) \right) 
\left( \mathrm{Ei}\left(\frac{\lambda h + \lambda_i}{\lambda g + \lambda_i}\right) - \mathrm{Ei}\left(\frac{\lambda_i}{\lambda g + \lambda_i}\right)\right) + 
\lambda_g \lambda_h e^{\frac{\lambda h (\lambda g \lambda_h + \lambda_h l)}{\lambda g + \lambda_i}} \right],
\]
where \( l > 0 \) and \( \mathrm{Ei}(x) = \int_{-\infty}^{x} e^t/t \, dt \). Hence, the outage probability is given as
\[
P(L < \gamma_h) = \int_{0}^{\gamma_h} f_L(l) \, dl.
\]
Consequently, based on the PDF of \( L \) in (47), the outage probability can be derived as (26).

**APPENDIX B: PROOF OF PROPOSITION 1**

Based on (1) and (25), the rate outage probability for the perfect timing synchronization, perfect SIC, and perfect CSI can be expressed as
\[
\hat{P}_{\text{out}} = P[\gamma_k < 2^{Y} - 1]
= \int_{0}^{\gamma_h} f_L(l) \, dl.
\]
As described in Section IV, \(|h_k|^2, |g_{k,i-1}|^2, |g_{k,i-2}|^2\) follow the exponential distributions with parameters \( \lambda_h, \lambda_g 1, \) and \( \lambda_g 2, \) respectively. Therefore, we can rewrite (49) in terms of SNR as
\[
\hat{P}_{\text{out}} = P[\hat{X} + \hat{Y} < g(\text{SNR})],
\]
where \( \hat{X} = \frac{|h_k|^2 p_k}{\frac{|h_k|^2 \sum_{i=1}^{N_k} p_i + \sigma^2}} \) and \( \hat{Y} = \frac{|g_{k,i-1}|^2 + |g_{k,i-2}|^2}{\sum_{i=1}^{N_k} p_i + \sigma^2} \).

Because of the independent channel gains, \( \hat{P}[\hat{X} + \hat{Y} < g(\text{SNR})] \leq \int_{0}^{\gamma_h} f_L(l) \, dl \), at high SNR, \( \lambda_h (2^{2Y-1}) \), we first have
\[
P[\hat{X} < g(\text{SNR})] = 1 - \exp \left[ - \frac{\lambda_h g(\text{SNR})}{\Phi_k - \sum_{i=1}^{N_k} (2^{2Y-1})} \right]
= 1 - \exp \left[ - \frac{\lambda_h (2^{2Y-1})}{\Phi_k - \sum_{i=1}^{N_k} (2^{2Y-1})} \right]
\sim \frac{\lambda_h}{\Phi_k - \sum_{i=1}^{N_k} (2^{2Y-1})} \cdot \left( \frac{2^{2Y-1}}{\Phi_k} \right).
\]

In addition, based on Fact-1 and Fact-2 in Appendix-I and Eq. (26) of [32] to find the outage behavior when \( \text{SNR} \to \infty \), we obtain
\[
P[\hat{Y} < g(\text{SNR})] \sim 2 \lambda_g^2 \left( \frac{2 Y - 1}{\text{SNR}} \right)^2.
\]
By substituting (51) and (52) into (50), we have the asymptotic rate outage probability in (27).

**APPENDIX C: PROOF OF LEMMA 2**

For the imperfect SIC, perfect timing synchronization and perfect CSI, we can rewrite (21) as \( L = A + B \). Let \( A \) and \( B \) be the RVs as defined in (41) and (43). Also, suppose \( F \sim \exp(\lambda_q) \) and \( \varpi = F + B \). Then, the PDF of \( \varpi \) is given as
\[
f_{\varpi}(\varpi) = \lambda_q \prod_{i=1}^{l} \lambda_i \left( \sum_{i=1}^{l} \left( \frac{e^{-\lambda_i \varpi}}{\lambda_i} \right) \right),
\]
where \( \varpi > 0 \). The PDF of \( Q = \frac{\varpi}{\lambda_q} \) is given as
\[
f_{Q}(q) = \frac{\psi_{\varpi}^{q} \psi_{\lambda}^{2} + q \lambda h \left( \lambda_q \sum_{i=1}^{l} (i + 2) q \lambda_i \psi_{\lambda}^{i-1} \right)}{(\lambda h \varpi + \lambda_q)^2 \prod_{i=1}^{l} (\lambda h \varpi + \lambda_q)^2)},
\]
where \( q > 0 \). Thus, the mean and variance of \( Q_{\varpi} \) are given by \( \mathbb{E}[Q] = \frac{\psi_{\varpi}^{q} \psi_{\lambda}^{2}}{\lambda_q} \) and \( \text{Var}[Q] = \frac{\psi_{\varpi}^{q} (\psi_{\lambda}^{2} + 2 q \lambda \psi_{\lambda} + 2 q^{2} \lambda^{2} \psi_{\lambda}^{2})}{\lambda_q^2} \), respectively. By convolving (51) and (52) into (54), we obtain the PDF (30) of \( Q_{\varpi} = \frac{\lambda_h \varpi}{\lambda_q + \lambda_h} \) is given as
\[
f_{Q_{\varpi}}(q) = \int_{0}^{\infty} f_{\lambda_q}(q) f_{\varpi}(\varpi) d\varpi.
\]
By substituting (41) and (55) into (56), we obtain the PDF (30) of \( Q_{\varpi} = \frac{\lambda_h \varpi}{\lambda_q + \lambda_h} \) given as
\[
f_{Q}(q) = \int_{0}^{\infty} f_{\lambda_q}(q) f_{\varpi}(\varpi) d\varpi.
\]
Table IV: K-S Test Results for Lemma 3

| Distms.  | Mean and Variance | Estimated Para. | MSE |
|----------|------------------|-----------------|-----|
|          |  $\mu$ | $\sigma$ | $\kappa_1$ | $\kappa_2$ | $\epsilon_1$ | $\epsilon_2$ |
| Gamma    | 5.9834 | 23.87 | 0.8059 | 0.3089 | 0.0016 | 0.0192 |
| Weibull  | 6.0164 | 24.9705 | 6.4089 | 1.2095 | 0.0177 | 0.0268 |
| Exponential | 5.9834 | 35.8037 | 5.9834 | 0.0189 |
| Rayleigh | 7.2086 | 14.1672 | 5.7452 | 0.0092 |
| Rician   | 7.2010 | 14.1687 | 0.1753 | 5.7442 | 0.3146 | 0.0102 |
| Nakagami | 6.4162 | 24.8484 | 0.4744 | 66.016 | 0.0017 | 0.3030 |

APPENDIX D: PROOF OF PROPOSITION 2

Based on (21) and (25), the rate outage probability for the perfect timing synchronization, imperfect SIC, and perfect CSI can be re-written as

$$P_{out}^w = P[\gamma_k^w < 2T - 1]$$

$$= P\left[1 - \frac{|h_k|^2\Phi_k}{\text{SNR}(\eta)|g_n|^2\Phi_n + \sum_{i=1}^{k-1} |h_i|^2\Phi_i} + 1 + \frac{|g_{k,k-1}|^2 + |g_{k,k-2}|^2}{2} < \frac{2T - 1}{\text{SNR}} \right].$$

(57)

If $\hat{Z} = \frac{|h_k|^2\Phi_k}{\text{SNR}(\eta)|g_n|^2\Phi_n + \sum_{i=1}^{k-1} |h_i|^2\Phi_i} + 1$, using power series expansion, we have

$$P[\hat{Z} < g(\text{SNR})] = 1 - \exp\left[-\hat{g}(\text{SNR})\right] \sim \hat{g}(\text{SNR}).$$

(58)

By substituting (58) and (52) into (57), we can approximate the outage behaviour when SNR $\rightarrow \infty$ as in (29). $\square$

APPENDIX E: PROOF OF LEMMA 3

For the perfect SIC, imperfect timing synchronization, and perfect CSI, we can rewrite (15) as $V = Q_1 + R$. Further, the second part of (13) can be expressed as $R = \chi$, where $\nu = \frac{(C + \epsilon_1^2)D}{C + D}$ and $\Lambda = \frac{\text{SNR}|\nu|^2}{C + D}$ with $|\nu|^2 = |1 - \epsilon_1|g_{4,1}g_{1,4} - \epsilon_2g_{4,1}g_{2,1}]^2$. It can be shown from (33) that $\nu$ is the Generalized Gamma distribution. Whereas, let $\Lambda$ be a Gamma distribution. Then, the PDF of $\Lambda$ is the Gamma distribution (33). Also, $R$ is a ratio of Generalized Gamma Distribution and Gamma Distribution (33). Thus, the PDF of $R$ is given as in (59). Similarly, suppose $\nu \sim \text{Exp}(\lambda_\nu)$, $D \sim \text{Exp}(\lambda_D)$, $0 < \epsilon_1 \leq 1$, and $0 < \epsilon_2 \leq 1$. Then, the PDF of $R = \frac{1}{1 - \epsilon_1}C + \epsilon_2C + D + \epsilon_2D$ is

$$f_R(r) = \frac{1}{4(\epsilon_1 - 1)}\lambda_\nu \left[1 \sqrt{\epsilon_2\gamma} \right] \left[\sqrt{\pi}erf\left(\sqrt{\lambda_\nu r}\right) - erf\left(\frac{\sqrt{\lambda_\nu r}}{\epsilon_1}\right)\right] + 2e^{-\frac{\lambda_\nu r}{\epsilon_1}} - 2e^{-\lambda_\nu(-r)},$$

(59)

where $r > 0$. Also, using this PDF, we can obtain the mean and variance as

$$\mathbb{E}[R] = \frac{2(1 + \epsilon_1 + \epsilon_1^2)}{3\lambda_\nu},$$

(60)

and

$$\text{Var}[R] = \frac{2(17 + 7\epsilon_1 - 3\epsilon_1^2 + 7\epsilon_1^3 + 17\epsilon_1^4)}{45\lambda_\nu^2},$$

(61)

respectively. By applying the Kolmogorov–Smirnov test (K-S test) [35, 36] on the distribution of $V = Q_1 + R$, it is determined that $V \sim \Gamma(\alpha, \beta)$. Further, the PDF of $V = Q_1 + R$ is given as

$$f_V(v) = \frac{\beta - \alpha e^{\alpha - \beta}}{\Gamma(\alpha)} v^\beta e^{-\beta v},$$

(62)

where $\alpha = \frac{\mathbb{E}[V]^2}{\text{Var}[V]}$ and $\beta = \frac{\text{Var}[V]}{\mathbb{E}[V]}$ are the parameters of Gamma distribution. Also, $\Gamma(x)$ is the gamma function given as $\Gamma(x) = (x - 1)!$. By applying the mathematical operation (30) and (37), as given in (48) on the PDF of $V$ in (62), we obtain the outage probability as shown in (30). $\square$

APPENDIX F: PROOF OF PROPOSITION 3

Based on (15) and (25), the rate outage probability for the imperfect timing synchronization, perfect SIC, and perfect CSI can be expressed as

$$P_{out}^e = P[\gamma_k^e < 2T - 1]$$

$$= P\left[1 - \frac{|h_k|^2\Phi_k}{\text{SNR}\sum_{i=1}^{k-1} |h_i|^2\Phi_i} + 1 + \frac{|g_{k,k-1}|^2 + |g_{k,k-2}|^2}{2} < \frac{2T - 1}{\text{SNR}} \right],$$

(63)

where $\nu_1 = |g_{k,k-1}|^2$, $\nu_2 = |g_{k,k-2}|^2$, and $|\nu|^2 = |(1 - \epsilon_1)g_{k,k-1}g_{k,k-2}g_{k,k-1} - \epsilon_2g_{k,k-2}g_{k,k-1}|^2$. Letting $W = \frac{|\nu|^2 + |\nu|^2}{2}$, based on Fact-1 and Fact-2 in Appendix-I of [32], we can obtain

$$P[|W| < g(\text{SNR})] \sim g^*(\text{SNR}),$$

(64)

where $g^*(\text{SNR}) = 2\epsilon_1\lambda_\nu^2 \left(\frac{2\epsilon_1}{\text{SNR}}\right)^\sigma$. By substituting (64) and (51) into (63), we can approximate the outage behaviour when SNR $\rightarrow \infty$ as in (31). $\square$

APPENDIX G: PROOF OF LEMMA 4

For the imperfect SIC, imperfect timing synchronization, and perfect CSI, we first rewrite (17) as $V_\eta = Q_\eta + R$, where the PDFs of $Q_\eta$ and $R$ are given in (35) and (59), respectively. Applying the Kolmogorov–Smirnov test (K-S test) [35, 36] on the distribution of $V_\eta = Q_\eta + R$, it is determined that $V_\eta \sim \Gamma(\theta, \phi)$, which corresponds to the following PDF

$$f_{V_\eta}(v) = \frac{\phi - \theta(v)\theta^{-1}e^{-\theta(v)}}{\Gamma(\theta)}(\theta(v)^\theta - 1),$$

(65)

where $v > 0$. In addition, $\theta = \frac{(\mathbb{E}[V_\eta])^2}{\text{Var}[V_\eta]}$ and $\phi = \frac{\text{Var}[V_\eta]}{\mathbb{E}[V_\eta]}$ are the parameters of Gamma distribution, where the mean and the variance of the $V_\eta$ are given by

$$\mathbb{E}[V_\eta] = \int_0^\infty v f_{V_\eta}(v) dv = \frac{\psi_1^2\psi_3}{\lambda_\nu^2} + 2(\epsilon_1^2 + \epsilon_1 + 1) \frac{3\lambda_\nu}{\lambda_\nu},$$

(66)
and
\[
\text{Var}[V_o] = \int_0^\infty v^2 f_{V_o}(v) dv - \mathbb{E}[V_o]^2
\]
\[
= \frac{\psi_0^2(\psi_0^2 + 2\psi_0^2)}{\psi_0^2} + 2 \left( 17\epsilon_1 + 7\epsilon_2 - 3\epsilon_1^2 + 7\epsilon_1 + 17 \right) \frac{45\lambda_b^2}{45^2},
\]
(67)
respectively. By substituting (65) into (48), as in (30), we can derive the outage probability as shown in (32).

**APPENDIX H: PROOF OF PROPOSITION 4**

With (17) and (25), the rate outage probability for the imperfect timing synchronization, imperfect SIC, and perfect CSI can be found as (17) is less than \( \gamma_{th} \), which can be written as
\[
\hat{P}_{out} = \mathbb{P}[\gamma_{out}^c < 2^T - 1]
\]
\[
= \mathbb{P}\left[ \frac{|h_k|^2 \Phi_k}{SNR} + \sum_{i=1}^{k-1} |h_i|^2 \Phi_i + 1
\right. \\
+ \left. \frac{(|\varphi|^2 + \epsilon_1|\varphi|^2)^2}{SNR(|\varphi|^2 + 2(|\varphi|^2 + |\varphi|^2)^2)} < 2^T - 1 \right].
\]
(68)
By substituting (52) and (64) into (68), we can obtain (33).

**APPENDIX I: PROOF OF LEMMA 5**

For the perfect SIC, perfect timing synchronization, and imperfect CSI given in (19), let \( A_p = \theta_{n-n-3}, \quad A_g = \theta_{n-n-3}, \quad B_p = \theta_{n-n-2}, \quad B_g = \theta_{n-n-2}, \quad C_\chi = |A|^2 |\varphi|^2 \), and \( D_\chi = |B_\chi|^2 |\varphi|^2 \), where \( A_p = A_g + \rho A_g \) and \( B_p = B_g + \rho B_g \). As described in Section III-C, \( A_p \sim CN(0, \sigma^2), \quad A_g \sim CN(0, \sigma^2), \quad B_p \sim CN(0, \sigma^2), \quad B_g \sim CN(0, \sigma^2) \). Therefore, we can find out that \( A_\chi \sim CN(0, \sigma^2 + \rho^2 \sigma^2) \) and \( B_\chi \sim CN(0, \sigma^2 + \rho^2 \sigma^2) \). We model \( A_\chi \) and \( B_\chi \) as mutually independent complex Gaussian RVs with variance \( \sigma^2 \). Then, their magnitudes (i.e., \( |A_\chi| \) and \( |B_\chi| \)) follow the Rayleigh distribution, and their squared magnitudes (i.e., \( |A_\chi|^2 \) and \( |B_\chi|^2 \)) follow the exponential distributions with the parameter \( \lambda_\chi \), where \( \lambda_\chi = 1/\sigma^2 = 1/(\sigma^2 + \rho^2 \sigma^2) \) (21), (22). The PDFs of \( C_\chi \) and \( D_\chi \) are given as \( f_{C_\chi}(c) = \frac{e^{-c/\lambda_\chi}}{\lambda_\chi} \), \( c > 0 \), and \( f_{D_\chi}(d) = \frac{e^{-d/\lambda_\chi}}{\lambda_\chi} \), \( d > 0 \), respectively. If we define another variable \( Z_\chi = C_\chi + D_\chi \), it follows a Gamma distribution, and its PDF is obtained by convolving \( f_{C_\chi}(c) \) and \( f_{D_\chi}(d) \) as
\[
f_{Z_\chi}(z) = \frac{z e^{-z/\lambda_\chi}}{\lambda_\chi^2}, \quad z > 0.
\]
(69)
For the perfect SIC, perfect timings and imperfect CSI, (19) can be written as \( L_\chi = Q_1 + Z_\chi \) and the PDF of \( L_\chi \) is obtained by convolving (45) and (69), which is given as
\[
f_{L_\chi}(l) = \frac{1}{\lambda_\chi^2} \sum_{i=1}^{I} \left[ \frac{\lambda_\chi \lambda_i}{\lambda_\chi \lambda_i} \right] \left( \lambda_i(-\lambda_\chi \lambda_i + \lambda_i + \lambda_i l) \right) \]
\[
\left( \frac{\lambda_i}{\lambda_\chi \lambda_i} \right) - \frac{\lambda_i}{\lambda_\chi \lambda_i} + \lambda_i l \right), \quad l > 0.
\]
(70)
As a result, the corresponding outage probability is obtained as (34).

**APPENDIX J: PROOF OF PROPOSITION 5**

From (19) and (25), the rate outage probability for the perfect timing synchronization, perfect SIC, and imperfect CSI can be expressed as
\[
\hat{P}_{out} = \mathbb{P}[\gamma_{out}^c < 2^T - 1]
\]
\[
= \mathbb{P}\left[ \frac{|h_k|^2 \Phi_k}{SNR} + \frac{(|A|^2 + |B_\chi|^2)}{2} < \frac{2^T - 1}{SNR} \right].
\]
(71)
Following the same procedure as in (52), we obtain the PDF of \( Z_\chi \) as
\[
\mathbb{P}[Z_\chi < g(SNR)] \sim 2\lambda_\chi^2 \left( \frac{2^T - 1}{SNR} \right)^2.
\]
(72)
Replacing (52) and (72) into (71), we have the asymptotic rate outage probability in (35).
H. Mheidat and M. Uysal, “Non-coherent and mismatched-coherent receivers for distributed STBCs with amplify-and-forward relaying,” *IEEE Trans. Wireless Commun.*, vol. 6, no. 11, pp. 4060–4070, 2007.

L. Lv, J. Chen, and Q. Ni, “Cooperative non-orthogonal multiple access in cognitive radio,” *IEEE Commun. Lett.*, vol. 20, no. 10, pp. 2059–2062, 2016.

Z. Zhang, Z. Ma, M. Xiao, Z. Ding, and P. Fan, “Full-duplex device-to-device-aided cooperative nonorthogonal multiple access,” *IEEE Trans. Vehicular Tech.*, vol. 66, no. 5, pp. 4467–4471, 2016.

L. Lv, J. Chen, Q. Ni, and Z. Ding, “Design of cooperative non-orthogonal multicast cognitive multiple access for 5G systems: User scheduling and performance analysis,” *IEEE Trans. Commun.*, vol. 65, no. 6, pp. 2641–2656, 2017.

M. Xu, F. Ji, M. Wen, and W. Duan, “Novel receiver design for the cooperative relaying system with non-orthogonal multiple access,” *IEEE Commun. Lett.*, vol. 20, no. 8, pp. 1679–1682, Aug. 2016.

R. E. Walpole, R. H. Myers, S. L. Myers, and K. Ye, *Probability and statistics for engineers and scientists*. Macmillan New York, 1993, vol. 5.

I. S. Gradshteyn and I. M. Ryzhik, *Table of integrals, series, and products*. Academic press, 2014.

J. N. Laneman, D. N. Tse, and G. W. Wornell, “Cooperative diversity in wireless networks: Efficient protocols and outage behavior,” *IEEE Trans. Info. Theory*, vol. 50, no. 12, pp. 3062–3080, 2004.

C. A. Coelho and J. T. Mexia, “On the distribution of the product and ratio of independent generalized gamma-ratio random variables,” *Sankhya: The Indian Journal of Statistics (2003-2007)*, vol. 69, no. 2, pp. 221–255, 2007. [Online]. Available: [http://www.jstor.org/stable/25664553](http://www.jstor.org/stable/25664553)

N. Wiener, *The Fourier integral and certain of its applications*. CUP Archive, 1988.

I. M. Chakravarti, R. G. Laha, and J. Roy, “Handbook of methods of applied statistics,” *Wiley Series in Probability and Mathematical Statistics (USA)* eng, 1967.

A. G. Glen, L. M. Leemis, and D. R. Barr, “Order statistics in goodness-of-fit testing,” *IEEE Trans. Reliability*, vol. 50, no. 2, pp. 209–213, Jun. 2001.

C. Bettstetter, H. Hartenstein, and X. Pérez-Costa, “Stochastic properties of the random waypoint mobility model,” *Wireless Networks*, vol. 10, no. 5, pp. 555–567, Sep 2004. [Online]. Available: [https://doi.org/10.1023/B:WINE.0000036458.88990.c5](https://doi.org/10.1023/B:WINE.0000036458.88990.c5)