DG2: A Faster and More Accurate Differential Grouping for Large-Scale Black-Box Optimization

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Abstract—Identification of variable interaction is essential for an efficient implementation of a divide-and-conquer algorithm for large-scale black-box optimization. In this paper, we propose an improved variant of the differential grouping algorithm, which has a better efficiency and grouping accuracy. The proposed algorithm, DG2, finds a reliable threshold value by estimating the magnitude of roundoff errors. With respect to efficiency, DG2 reuses the sample points that are generated for detecting interactions and saves up to half of the computational resources on fully separable functions. We mathematically show that the new sampling technique achieves the lower bound with respect to the number of function evaluations. Unlike its predecessor, DG2 checks all possible pairs of variables for interactions and has the capacity to identify overlapping components of an objective function. On the accuracy aspect, DG2 outperforms the state-of-the-art decomposition methods on the latest large-scale continuous optimization benchmark suites. DG2 also performs reliably in the presence of imbalance among contribution of components in an objective function. Another major advantage of DG2 is the automatic calculation of its threshold parameter ($\epsilon$), which makes it parameter-free. Finally, the experimental results show that when DG2 is used within a cooperative co-evolutionary framework, it can generate competitive results as compared to several state-of-the-art algorithms.

Index Terms—large-scale global optimization, problem decomposition, differential grouping, cooperative co-evolution.

I. INTRODUCTION

Large-scale global optimization has become an active field of research in the past decade due to the growing number of large-scale optimization problems in engineering and sciences [1, 2]. Most engineering problems have shown an exponential increase in the number decision variables they entail [3]. Advances in machine learning and the rise of deep artificial neural networks has resulted in optimization problems with over a billion variables [4, 5]. Ubiquity of data has also caused the emergence of large-scale optimization problems at the heart of many data analytics and learning problems [6]. Target shape design optimization for aircraft wings and turbine blades [7], satellite layout design [8], parameter estimation in large scale systems biology models [9], seismic waveform inversion [10], and parameter calibration of water distribution system [11] are just a few examples from a wide array of large-scale optimization problems.

A major challenge of large-scale optimization is the exponential growth in the size of the search space with respect to the number of decision variables. It is this “curse-of-dimensionality” that has made large-scale optimization an exceedingly difficult task. This motivated the development of a wide range of scalable algorithms in the classic mathematical programming domain [12, 13] as well as metaheuristics [14, 15]. Evolutionary algorithms, in particular, have shown superior performance as compared to other classic methods on problems with millions or even billions of variables [16, 17]. Other methods such as swarm intelligence [18–20], memetic algorithms [21–23], differential evolution [24, 25], evolution strategies [26], and estimation of distribution algorithms [27, 28] have also gained popularity for large-scale optimization because of their ability to deal with black-box problems. It should be noted that the notation of large-scale changes over time and varies from problem to problem. In a broad sense, a problem is considered large-scale if it causes scalability issues on the state-of-the-art algorithms. For the current study, which focuses on real-parameter optimization, the existing algorithms exhibit scalability issues on problems having more than about a hundred decision variables.

A number of approaches such as dimensionality reduction [29], surrogate modelling [30], local search [21, 22], and divide-and-conquer (a.k.a decomposition) methods [31] can be used for large-scale optimization, among which decomposition methods have gained popularity in recent years [18, 31–47]. Decomposition methods break a large-scale problem into a set of smaller and simpler subproblems each of which is optimized in an iterative manner. In the context of evolutionary algorithms, cooperative co-evolution [48] is a popular means of exploiting the modular nature of many complex large-scale problems and has been used in a wide range of areas beyond optimization such as the study of evolutionary game theory [49]. A major challenge of using a cooperative co-evolutionary framework for large-scale optimization is the right choice of problem decomposition. Ideally, a given objective function should be decomposed such that the interaction between the resultant components is minimized. For a black-
box optimization problem, the variable interaction information are not available. Therefore, specific algorithms are required to identify the underlying interaction structure of the decision variables.

Differential Grouping (DG) [32] is a competitive decomposition algorithm that can identify the nonseparable components of a continuous objective function and has shown superior performance as compared to other decomposition algorithms such as variable interaction learning [33] on the CEC’2010 [50] large-scale benchmark suite [32]. Despite its success on the CEC’2010 benchmark problems, it has been shown that differential grouping has some difficulty with the CEC’2013 large-scale benchmark functions [51]. In particular, differential grouping has the following major shortcomings:

- High computational cost on fully separable functions.
- Inability to detect objective functions with overlapping components, i.e., components that share decision variables [51].
- Sensitivity to computational roundoff errors [37].
- Requiring the user to specify a threshold parameter ($\epsilon$).

In this paper, we propose an improved version of differential grouping that addresses the above issues. In particular, this improved version, DG2, reduces the total number of objective function evaluations by half for fully separable functions which require the most function evaluations. This allows the algorithm to check all pairs of variables for interaction at a much lower cost as compared to its predecessor. Testing all pairs of variables for interaction is essential to identify functions with overlapping components. The reduction in the total number of objective function evaluations is achieved through systematic generation of sample points to maximize point reuse in the process of applying the differential grouping theorem (see Section II). We mathematically show that this new method achieves the lower bound when the differential grouping theorem is used to detect the interactions.

In addition to improving the efficiency, DG2 significantly improves the grouping accuracy of differential grouping on the existing large-scale benchmark suites. A major advantage of DG2 is its parameter-free property. DG2 takes the computational rounding errors into account in estimating a proper threshold value ($\epsilon$) which determines its sensitivity to weak interactions. In particular, DG2 has the following advantages over the static method used in differential grouping:

- Unlike DG that uses a single global $\epsilon$ value to detect all the interactions, DG2 dynamically calculates an $\epsilon$ value to detect the interaction between each pair of variables. For each interaction, DG2 approximates the magnitude of roundoff errors and calculates the threshold value accordingly. This is particularly useful when dealing with imbalanced functions, in which the magnitude of roundoff error may be different from component to component.
- Unlike DG, the new method does not require the user to specify any external parameter. In other words, DG2 is parameter-free.

The organization of the rest of this paper is as follows. Section II contains the details of the differential grouping theorem and algorithm. Section III gives an outline of the proposed improvements. Section III-A contains the details on how to reduce the total required objective function evaluations as well as a proof of the lower bound for the total required evaluations. Section III-B focuses on improving the accuracy of differential grouping and making it parameter-free. The experimental results about the grouping accuracy of DG2 and its performance within a cooperative co-evolutionary framework are presented in Section IV. Finally, Section V concludes the paper.

II. BACKGROUND AND RELATED WORK

Problem decomposition is an integral part of using cooperative co-evolution for function optimization. A good problem decomposition is one that has minimal dependence among its components. This is often characterized by separability structure of the objective function, which is defined as follows:

**Definition 1** ([51]). A function $f(x)$ is partially separable with $m$ independent components iff:

$$\arg \min_{x} f(x) = \left( \arg \min_{x_1} f(x_1, \ldots, x_m), \ldots, \arg \min_{x_m} f(\ldots, x_m) \right),$$

where $x = (x_1, \ldots, x_n)^\top$ is a decision vector of $n$ dimensions, $x_1, \ldots, x_m$ are disjoint sub-vectors of $x$, and $2 \leq m \leq n$.

Additive separability is a special type of partial separability, which is defined as follows:

**Definition 2** ([51]). A function is partially additively separable if it has the following general form:

$$f(x) = \sum_{i=1}^{m} f_i(x_i), \quad m > 1,$$

where $f_i(\cdot)$ is a non-separable subfunction, and $m$ is the number of non-separable components of $f$. The definition of $x$ and $x_i$ is identical to what was given in Definition 1.

Many decomposition algorithms have been proposed to decompose a black-box optimization problems into smaller subproblems. Static grouping is the simplest decomposition strategy in which the decision variables are grouped into arbitrary groups. In its simplest form, an $n$-dimensional problem is broken down into $s k$-dimensional problems. Examples of such methods are the divide-in-half method by Shi et al. [52], and the method employed by van den Bergh and Engelbrecht [53]. These methods are oblivious of variable interactions which may have a significant impact on the optimization performance [31]. Some other decomposition algorithms such as random grouping [31], adaptive variable partitioning [39], delta grouping [54], and min/max variance decomposition [55] use various heuristics in order to form the groups based on variable interaction characteristics of the objective function. The drawback of these methods is their low grouping accuracy, and the fact that they presuppose the number and/or the size of components. These algorithms also divide the decision variables into $s k$-dimensional components. Improved versions of random grouping and delta grouping use a so-called multi-level strategy [54, 56] in which multiple fixed decompositions are used over the course of optimization. More sophisticated decomposition methods such as variable interaction
learning [33], meta-modelling decomposition [40], statistical learning decomposition [38], and differential grouping [32] do not presuppose the number and/or size of components. Among these algorithms, differential grouping has shown superior performance with respect to grouping accuracy [33, 40]. The following theorem is at the heart of interaction detection of differential grouping:

**Theorem 1** ([32]). Let \( f(x) \) be an additively separable function. \( \forall a, b_1 \neq b_2, \delta \in \mathbb{R}^1, \delta \neq 0 \), variables \( x_p \) and \( x_q \) interact if the following condition holds

\[
\Delta_{\delta, x_p}[f](x)|_{x_p=a, x_q=b_1} \neq \Delta_{\delta, x_p}[f](x)|_{x_p=a, x_q=b_2},
\]

where

\[
\Delta_{\delta, x_p}[f](x) = f(..., x_p + \delta, ...) - f(..., x_p, ...),
\]

refers to the forward difference of \( f \) with respect to variable \( x_p \) with interval \( \delta \).

Theorem 1 states that two variables \( x_p \) and \( x_q \) interact if Equation (2) evaluated with any two different values of \( x_q \) gives different results [32]. A proof of this theorem can be found in [32]. For the sake of brevity the left hand side of Equation (1) is denoted by \( \Delta^{(1)} \) and its right hand side by \( \Delta^{(2)} \). It is clear that \( \Delta^{(1)} \neq \Delta^{(2)} \iff |\Delta^{(1)} - \Delta^{(2)}| \neq 0 \). It is also clear that this equality check is not practical on computing devices due to limited precision of floating-point numbers. For this reason, the equality check can be converted into an inequality check of the form \( \lambda = |\Delta^{(1)} - \Delta^{(2)}| > \epsilon \) by introducing the control parameter \( \epsilon \) that determines the sensitivity of differential grouping to interactions.

Two major drawbacks of differential grouping are its sensitivity to the parameter \( \epsilon \) and its poor accuracy in detecting interacting variables on functions with overlapping components. As reported in [32], the grouping accuracy of differential grouping is low on the Rosenbrock function [57] which has overlapping components with overlap size of one. Also, if differential grouping is used to find the interaction structure of functions with overlapping variables, the shared decision variables between two components will be placed in one group and will be excluded from other groups. It is not yet clear what an optimal decomposition may be for an overlapping function; nevertheless, an accurate identification of the underlying structure is essential to propose a meaningful decomposition.

Global Differential Grouping (GDG) [37] and eXtended Differential Grouping (XDG) [58] are two variants of differential grouping, which aim at addressing the above shortcomings. XDG focuses on identifying indirect interactions in order to deal with the Rosenbrock function. The issue with XDG is that it inherits the sensitivity issue of differential grouping and also its method of inferring variable interaction may consider separable variables as nonseparable. This issue is discussed further in Section IV-A. GDG addresses the sensitivity issue of differential grouping by taking computational errors into account. However, the use of a global parameter to detect all interactions makes it unsuitable for imbalanced functions. This issue is discussed further in Section IV-B. GDG also addresses the problem of identifying overlapping functions by examining all pairs of variables for interaction. However, we will show in Section IV that DG2 can achieve the same goal with fewer objective function evaluations.

### III. IMPROVED DIFFERENTIAL GROUPING

In this section, we describe the details of improving the grouping accuracy and efficiency of differential grouping. Algorithm 1 shows the high-level structure of DG2 that incorporates these improvements. DG2 has three major parts. The first part is forming what we call a raw interaction structure matrix (\( \Lambda \)) that contains the quantity \( |\Delta^{(1)} - \Delta^{(2)}| \) for all pairs of variables. This is done by the ISM function. The second part of the algorithm is finding a suitable threshold parameter (\( \epsilon \)) in order to convert the raw interaction structure matrix \( \Lambda \) to a design structure matrix \( \Theta \). The entry \( \Theta_{ij} \) takes 1 if \( \Lambda_{ij} > \epsilon \), and 0 otherwise. It should be noted that unlike differential grouping and GDG, DG2 obtains a threshold based on information such as magnitude of function values and the values of the raw interaction structure matrix that are calculated by the ISM function (Algorithm 1). Finally, the last part of the algorithm deals with the decomposition of the variables into nonseparable groups, which is performed by identifying the connected components of the graph with the node adjacency matrix \( \Theta \). This can be efficiently done in linear time in \( n \) [59].

It should be noted that a complete design structure matrix is necessary to detect overlapping functions, in which different components share common variables. This type of functions is more general in practice, and is more challenging. Given the design structure matrix, various decompositions can be devised in order to deal with overlapping components. However, the study of an optimal decomposition for overlapping functions is beyond the scope of this study.

In the remainder of this paper, we focus on two major issues:

- Finding an efficient implementation for the ISM function in order to form the interaction structure matrix using the minimum possible function evaluations (Section III-A).
- Finding an effective thresholding method that results in an accurate decomposition of a function into its components that generalizes over a wide range of functions (Section III-B).

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1 Values of \( a, b_1, b_2 \) and \( \delta \) are chosen such that \( f \) is evaluated within its domain.

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**Algorithm 1:** (\( g, x_1, ..., x_q, x_{sep}, \Gamma \) = DG2(\( f, n, \mathbb{X}, \mathbb{X} \))

1. \( (A, F, f, f_{base}, \Gamma) = ISM(f, n, \mathbb{X}, \mathbb{X}); \)
2. \( \Theta = DSM(A, F, f, f_{base}, n); \)
3. \( (k, y_1, ..., y_k) = ConnComp(\Theta); \)
4. \( x_{sep} = \{ \}, g = 0; \)
5. for \( i = 1 \rightarrow k \) do
   6.   if \( |y_i| = 1 \) then
      7.     \( x_{sep} = x_{sep} \cup y_i; \)
   8.   else
      9.     \( g = g + 1, x_q = y_i; \)

A. Improving the Efficiency of Differential Grouping

As mentioned earlier, in order to detect the overlapping functions, it is essential to examine all pairs of variables for interaction. It is clear that for an n-dimensional function, the total number of interactions is \( \binom{n}{2} \). According to Theorem 1 each comparison requires four fitness evaluations which results in a total of \( 4 \times \binom{n}{2} = 2n(n-1) \) evaluations.

In this section, we show that by systematic selection of sample points for calculating the difference equation (2), the total number of fitness evaluations can be significantly reduced. In order to show this, we assume a simple function with only three decision variables, i.e., \( f(x_1, x_2, x_3) \). The total number of function evaluations according to Theorem 1 is as follows:

\[
x_1 \leftrightarrow x_2; \Delta^{(1)} = f(a', b, c) - f(a, b, c) \]
\[
x_1 \leftrightarrow x_3; \Delta^{(1)} = f(a', b, c) - f(a, b, c) \]
\[
x_2 \leftrightarrow x_3; \Delta^{(1)} = f(a, b', c) - f(a, b, c) \]

where \( a, b, \) and \( c \) are the values taken by \( x_1, x_2, \) and \( x_3 \) respectively, and \( a' = x_1 + \delta, \ b' = x_2 + \delta, \) and \( c' = x_3 + \delta. \)

For a clearer illustration, the points that are evaluated with function \( f \) are color-coded and are shown geometrically in Figure 1. From previous calculations we know that the total number of function evaluations for a 3-dimensional function is \( 2n(n-1)|_{n=3} = 12 \). However, it is clear from Figure 1 that only 7 unique points are required.

![Fig. 1: Geometric representation of point generation in DG2 for a 3D function.](image)

In order to calculate \( \Delta^{(1)} \) and \( \Delta^{(2)} \), four points are required. According to Theorem 1 these points are chosen such that they form a rectangle. To calculate \( \Delta^{(1)} \) a base point is required which, in this example, is \( (a, b, c) \). Then, in order to find interactions with \( x_1 \), the first variable should be varied in order to calculate \( \Delta^{(1)} \). Therefore, the second point will be \( (a', b, c) \). To find the interaction between \( x_1 \) and \( x_2 \), the same difference equation as \( \Delta^{(1)} \) should be evaluated for a different value of \( x_2 \). Therefore, we get \( (a', b, c') \) and \( (a, b', c) \). If we follow this pattern to find all interactions, we can see that the base point \( (a, b, c) \) is repeated exactly three times, the cases where only one dimension is varied with respect to the base point such as \( (a', b, c), (a', b', c) \), and \( (a, b, c') \) are repeated exactly two times, and the cases where two dimensions are varied with respect to the base point such as \( (a', b', c), (a', b', c') \), and \( (a, b', c') \) are evaluated only once.

![Algorithm 2](image)

This process can be generalized for an arbitrary number of decision variables. For a general case, we need the following evaluations in order to detect the interaction between the \( i \)th and the \( j \)th dimensions.

\[
\begin{align*}
\Delta^{(1)} &= f(\ldots, x'_i, \ldots) - f(x_1, \ldots, x_n) \\
\Delta^{(2)} &= f(\ldots, x'_i, \ldots, x'_j, \ldots) - f(\ldots, x'_j, \ldots)
\end{align*}
\]

Based on this pattern we see that the total number of evaluations is \( 2n(n-1) \). It should be noted that the number of unique evaluations is much less than this quantity due to redundant evaluations caused by the assumptions made previously.

\[
\begin{align*}
\left( n(n-1) \right) / 2 - 1 & : \text{redundant evaluations of } (x_1, \ldots, x_n) \\
\left( n(n-2) \right) / 2 & : \text{redundant evaluations of } (\ldots, x'_j, \ldots)
\end{align*}
\]

Therefore, to calculate the total number of unique evaluations, the number of redundant evaluations should be subtracted from the total, which yields the following:

\[
\frac{n(n+1)}{2} + 1.
\]

In Theorem S.1, we show that this is the minimum number of objective function evaluations needed to form the interaction structure matrix. Theorem S.1 and its proof can be found in the supplementary document accompanying this paper (Section S-1). Algorithm 2 is an implementation of the process that was described above and achieves the lower bound according to Theorem S.1. The ISM function, generates the interaction structure matrix (\( \Lambda \)) which is used by the DSM function to find a reliable \( \epsilon \) to establish the separability or nonseparability of all pairs of variables.

B. Improving the Grouping Accuracy of Differential Grouping

It was mentioned in Section II that the grouping accuracy of differential grouping depends on \( \epsilon \). Theoretically, the value of

\[2\] A more detailed derivation is provided in Section S-1 of the supplementary document accompanying this paper.
\( \epsilon \) can be set to zero, since any positive difference between \( \Delta^{(1)} \) and \( \Delta^{(2)} \) implies an interaction between the variables in examination. However, in practice, the floating-point operations incur computational roundoff errors and cause nonzero \( \lambda \) values even for separable variables. A major challenge for differential grouping is to distinguish between a genuine nonzero \( \lambda \) due to variable interaction, and a nonzero \( \lambda \) due to computational errors. In this section, we show that the magnitude of roundoff errors is a function of the magnitude of the quantities used in a calculation. This makes a static threshold, such as the one used in differential grouping, an ineffective method. Omidvar et al. [51] have shown that the nonuniform contribution of components in an objective function significantly affects the accuracy of differential grouping with a static threshold. The functions in the CEC’2013 large-scale benchmark suite have such an imbalance property.

The new method of calculating a threshold value estimates the greatest lower bound (\( \epsilon_{\text{inf}} \)) and the least upper bound (\( \epsilon_{\text{sup}} \)) for the roundoff error by a mechanism which will be explained later. These values are calculated separately for each pair of variables based on the available information such as function values and the quantity \( \lambda = |\Delta^{(1)} - \Delta^{(2)}| \) to maximize the detection accuracy. Once the bounds are found, two variables are considered to interact if \( \lambda > \epsilon_{\text{sup}} \), and separable if \( \lambda < \epsilon_{\text{inf}} \). In order words, this interval defines a safe region that determines genuine zero and nonzero \( \lambda \) values. The \( \lambda \) values that fall outside this region may or may not be genuine nonzero values. To overcome this, we calculate the relative proportion of genuine zero and nonzero values and use it to bias the threshold towards either \( \epsilon_{\text{inf}} \) or \( \epsilon_{\text{sup}} \). The details of this process is given next.

Based on the IEEE 754 standard [60], the mapping of a real number \( x \) to a floating-point number (denoted by \( f(x) \)) may impose a rounding to the nearest representable number. According to the IEEE 754 standard, the representation error for a number \( x \) is a function of itself because: \( f(x) = x(1 + \delta) = x + \delta x \), where the bounds for \( \delta \) is determined by a machine dependent constant called the machine epsilon (\( \mu_{\text{M}} \)) such that \( |\delta| < \mu_{\text{M}} \) (see Theorem S.2 in Section S-II of the supplementary document). Therefore, the error term \( \delta x \) will grow with \( x \). Since differential grouping may deal with large numbers, it is essential to take the magnitude of the function values into account when estimating the threshold value (\( \epsilon \)).

In addition to the representational rounding error that was explained above, the floating-point arithmetic also incurs computational rounding error. The IEEE standard guarantees that \( x + y = f(x + y) \), where \( \oplus \) represents floating-point summation operator\(^3\). In other words, the floating-point sum of two numbers is guaranteed to be equal to the floating-point number closest to the real sum of the two numbers. In most models of error analysis, this is generalized to other operations such as subtraction, multiplication, division, and sometimes the square root function. This statement does not hold for a sequence of floating-point operations such as \( x_1 \oplus x_2 \oplus + \cdots \oplus x_n \) due to the accumulation of errors.

Theorem S.3 [61] can be used to find an upper bound for the accumulated arithmetic error in any calculation. An example of applying Theorem S.3 to calculate an upper bound for a dot-product example is given in Section S-II-A. In this paper, we use Theorem S.3 to find a reasonable upper and lower bounds for the error involved in calculating \( \Delta \). To estimate the greatest lower bound (infimum) for the magnitude of the roundoff error, we assume that the calculation of \( f(x) \) is error free, and the only source of error is in the application of differential grouping, i.e., the calculation of \( \lambda = |\Delta^{(1)} - \Delta^{(2)}| \).

Thus,
\[
\hat{\Delta}_1 = f(x) \oplus f(x') = (f(x) - f(x'))(1 + \delta_1) = \Delta^{(1)}(1 + \delta_1),
\]
\[
\hat{\Delta}_2 = f(y) \oplus f(y') = (f(y) - f(y'))(1 + \delta_2) = \Delta^{(2)}(1 + \delta_2),
\]
\[
\hat{\lambda} = |\hat{\Delta}_1 \ominus \hat{\Delta}_2| = |\hat{\Delta}_1 - \hat{\Delta}_2|(1 + \delta_3)
\]
\[
= |f(x)(1 + \delta_1)(1 + \delta_3) - f(x')(1 + \delta_1)(1 + \delta_4) - f(y)(1 + \delta_2)(1 + \delta_3) + f(y')(1 + \delta_2)(1 + \delta_3)|.
\]

We can see that the maximum number of products of the form \( (1 + \delta) \) is 2 \( (k = 2) \). Therefore, by applying Theorem S.3 we have:
\[
|\lambda - \hat{\lambda}| \leq \gamma_2 \left| (f(x) - f(x')) - (f(y) - f(y')) \right| \tag{3}
\]
\[
= \gamma_2 \left| (f(x) + f(y') - (f(y) + f(x')) \right|
\]
\[
\leq \gamma_2 \cdot \max \left\{ f(x) + f(y') \right\} := \epsilon_{\text{inf}}.
\]

Equation (3) is based on the assumption that the codomain of \( f \) is non-negative, i.e., \( f : \mathbb{R} \to \mathbb{R}_+^n \). A more general form for \( f : \mathbb{R} \to \mathbb{R} \) is as follows:
\[
\epsilon_{\text{inf}} = \gamma_2 \left( |f(x)| + |f(y')| + |f(y)| + |f(x')| \right) \tag{4}
\]

In this paper, the calculation of \( \epsilon_{\text{inf}} \) is based on Equation (3).

To estimate an upper bound for the roundoff error, we cannot assume that function evaluations are error free. However, the difficulty here is that the functions are black-box; therefore, we do not know the exact number of error terms \( (1 + \delta_i) \) in the calculation of \( f(\cdot) \). As a rule of thumb in the field of error analysis, it is customary to assume that the error grows with the square root of the number of floating-point operations \( \phi \) involved in a calculation [62]. In other words, to calculate an upper bound for the error based on Theorem S.3, we assume that \( k \approx \sqrt{\phi} \). We also assume that the error in calculating \( |\lambda - \hat{\lambda}| \) is negligible with respect to the error in \( f(\cdot) \). Therefore, an estimate of the least upper bound can be calculated as follows:
\[
|f(\cdot) - \hat{f}(\cdot)| \leq \gamma_\sqrt{\phi} f(\cdot) := \epsilon_{\text{sup}} \tag{5}
\]

The problem with Equation (5) is that in black-box optimization, we do not know the number of floating-point operations involved in calculating the objective function \( f(\cdot) \). To overcome this difficulty, instead of finding the exact number of floating-point operations, we make some assumptions about the relationship between the dimensionality of the problem \( (n) \) and the number of floating-point operations \( \phi \) that it may require. The simplest mapping of \( n \) variables into a scalar by a series of floating-point operations can be done by a simple summation, i.e., \( f(x) = \sum_{i=1}^{n} x_i \). In this example, the total

\(^3\)All other basic floating-point operations are shown in a circle in a similar way.
number of floating-point operations is \( \phi(n) = n - 1 \). The dot-product of two \( n \)-dimensional vectors is another common example, i.e., \( \mathbf{x} \cdot \mathbf{y} = \sum_{i=1}^{n} x_i y_i \). This calculation requires \( n \) multiplications and \( n-1 \) additions. Therefore, the total number of operations required to calculate the dot-product of two \( n \)-dimensional vectors is \( \phi(n) = 2n - 1 \). At a higher level of abstraction, we can say that the number of operations involved in the calculation of these functions is of order \( O(n) \). An example of a function with floating-point complexity of order \( O(n^2) \) is \( \sum_{i=1}^{n} x_i^2 \). It turns out that these complexity classes constitute a large body of numerical operations. Matrix operations are among the most computationally expensive, whose complexity does not exceed \( O(n^3) \). Table S-I contains a short list of common numerical operations and their complexity classes [62] (supplementary document Section S-III).

In an ideal situation, we need to find the least upper bound (\( e_{\text{sup}} \)) of the roundoff errors. If such a bound is available, any \( \lambda \) larger than \( e_{\text{sup}} \) can be treated as a genuine nonzero value. Generally speaking, the bounds calculated based on Theorem S.3 are very conservative, and the actual roundoff errors are much smaller in practice [61]. For example, Sterbenz’s Theorem [63] states that \( x + y = x - y \) if \( y/2 \leq x \leq 2y \). In other words, if two floating-point numbers are sufficiently close, their floating-point subtraction is exact. Additionally, modern computers have a fused multiply-add (FMA) instruction that involves a floating-point multiplication followed by an addition. Although FMA involves three floating-point operations, it commits only one rounding error in the worst case. We will also show in Section IV-D that underestimation of \( e_{\text{sup}} \) is not detrimental to detection of interacting variables. In general, underestimation of \( e_{\text{sup}} \), results in accurate detection of interacting variables at the expense of missing some separable variables. Conversely, overestimation of \( e_{\text{sup}} \), results in high detection accuracy of separable variables at the expense of missing interacting variables. It is clear that the latter case is more detrimental to the optimization performance. Therefore, to get a tighter bound, we assume a linear complexity and define \( e_{\text{sup}} \) as follows:

\[
e_{\text{sup}} = \gamma \sqrt{\pi} \max \{ f(x), f(x'), f(y), f(y') \}
\]

(6)

By estimating the least upper bound (\( e_{\text{sup}} \)) and the greatest lower bound (\( e_{\text{inf}} \)), we can identify reliable \( \lambda \) values. More specifically, all the \( \lambda \) values greater than \( e_{\text{sup}} \) will be treated as genuine nonzero (interacting variables), and all the values smaller than \( e_{\text{inf}} \) are treated as genuine zeros (separable variables). Finally, for the values in the range (\( e_{\text{inf}}, e_{\text{sup}} \)), the following weighted average of the bounds is used to set the threshold:

\[
\epsilon = \frac{\eta_0}{\eta_0 + \eta_1} e_{\text{inf}} + \frac{\eta_1}{\eta_0 + \eta_1} e_{\text{sup}},
\]

(7)

where \( \eta_0 \) is the number of entries in \( \Lambda \) which are less than \( e_{\text{inf}} \), and \( \eta_1 \) is the number of entries in \( \Lambda \) which are greater than \( e_{\text{inf}} \). Equation (7) is a natural choice for setting a threshold for non-reliable \( \lambda \) values. If \( \eta_0 = \eta_1 \), then Equation 7 reduces to the arithmetic mean of \( e_{\text{inf}} \) and \( e_{\text{sup}} \). This is intuitive, because when the number of reliably detected separable and nonseparable variables is equal, the middle point of the interval between \( e_{\text{inf}} \) and \( e_{\text{sup}} \) is the least biased choice of \( \epsilon \). Conversely, if the number of reliable calculations is skewed to one side, the threshold value should be biased to the same side. In the extreme case, if \( \eta_0 = 0 \), then \( \epsilon = e_{\text{sup}} \). Similarly, if \( \eta_1 = 0 \), then \( \epsilon = e_{\text{inf}} \).

Algorithm 3 contains the details of the DSM algorithm, in which the threshold on \( \Lambda \) is calculated by considering roundoff errors. The goal of Algorithm 3 is to convert the interaction structure matrix \( \Lambda \), which is calculated by DSM, into a binary design structure matrix (\( \Theta \)) that represent variable interactions.

### IV. Experimental Results and Analysis

In this section, we briefly compare the efficiency of DG2 with several state-of-the-art decomposition algorithms, namely Differential Grouping (DG) [32], GDG [37], XDG [58], and CCVIL [33]. Next, we assess the grouping accuracy of DG2 on the CEC’2010 and the CEC’2013 large-scale benchmark suites, and compare it with several other state-of-the-art decomposition algorithms. Next, we use DG2 in a cooperative co-evolutionary framework to test its efficiency on the final optimization performance. Finally, we assess the sensitivity of DG2 to the imbalance level between the components of an objective function, and the assumptions about complexity of floating-point operation in black-box functions.

#### A. Comparative Analysis of Grouping Efficiency

According to Theorem S.1, DG2 requires the least number of function evaluations to detect all interactions as compared to other decomposition algorithms. The total number of function evaluations needed by DG2 is constant and is equal to \( n^2 + n + 2 \) for an \( n \)-dimensional problem (Theorem S.1). The GDG algorithm requires \( n^2 + n + 2 \) evaluations [37], which is larger than what is needed by DG2 to construct the entire interaction structure matrix. Unlike DG2 and GDG, XDG
TABLE I: Grouping accuracy of DG2, XDG, GDG, DG, and CCVIL on the CEC’2010 and 2013 large-scale benchmarks. \( \rho_1 \) measures the accuracy of detecting interactions, \( \rho_2 \) measures the accuracy of detecting separable variables, and \( \rho_3 \) measures the overall accuracy. DG2 generalizes better on the CEC’2013 benchmarks and outperforms other algorithms by a wide margin.

| Benchmarks | Statistics | CCVIL | DG | GDG | XDG | DG2 |
|------------|------------|-------|----|-----|-----|-----|
|            |            | \( \epsilon = 10^{-3} \) | \( \epsilon = 10^{-3} \) | \( \epsilon = 10^{-6} \) | \( \alpha = 10^{-8} \) | \( \alpha = 10^{-10} \) | \( \epsilon = 10^{-3} \) |
| CEC’2010   | Median     | 3.5  | 100 | 99.7 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 |
|            | Std.       | 1.6  | 0.05 | 22.0 | 12.6 | 7.5 | 7.2 | 0.2 | 10.9 | 10.7 | 0.0 | 12.9 | 12.8 | 0.1 | 6.3 | 6.1 | 0.0 | 18.4 | 17.9 | 0.0 | 22.3 | 21.7 | 0.0 | 33.7 | 37.7 | 0.0 | 26.3 | 25.7 |
| CEC’2013   | Median     | 2.6  | 99.9 | 93.4 | 95.5 | 97.5 | 97.5 | 99.9 | 93.4 | 93.7 | 100 | 87.7 | 88.4 | 100 | 98.5 | 98.5 | 100 | 95.0 | 95.2 | 100 | 93.7 | 94.1 | 100 | 99.7 | 77.0 | 78.2 | 100 | 90.2 | 90.7 |
|            | Std.       | 1.0  | 0.05 | 22.0 | 12.6 | 7.5 | 7.2 | 0.2 | 10.9 | 10.7 | 0.0 | 12.9 | 12.8 | 0.1 | 6.3 | 6.1 | 0.0 | 18.4 | 17.9 | 0.0 | 22.3 | 21.7 | 0.0 | 33.7 | 37.7 | 0.0 | 26.3 | 25.7 |
| Success rate |          | 2.6  | 100 | 99.8 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 | 100 |
|            |            | 2.2  | 100 | 99.0 | 87.1 | 94.4 | 94.6 | 93.3 | 93.7 | 94.2 | 98.5 | 87.4 | 88.5 | 100 | 98.5 | 98.5 | 100 | 95.0 | 95.2 | 100 | 93.7 | 94.1 | 100 | 99.7 | 77.0 | 78.2 | 100 | 90.2 | 90.7 |
|            |            | 1.0  | 0.25 | 21.8 | 20.8 | 9.4 | 8.6 | 11.7 | 9.3 | 8.6 | 4.6 | 11.0 | 10.8 | 30.5 | 17.5 | 16.7 | 27.6 | 24.7 | 23.7 | 23.1 | 26.4 | 25.3 | 4.9 | 48.2 | 45.0 | 8.4 | 28.9 | 27.9 |

![Fig. 2: The interaction structures represented by (a) and (b) cannot be distinguished by XDG.](image)

does not construct a full interaction structure matrix and cannot identify the overlapping functions. If XDG detects that variables \( x_i \) and \( x_j \) both interact with a common variable \( x_k \), it does not check the interaction between \( x_i \) and \( x_j \) explicitly. Therefore, XDG fails to distinguish between the interaction structures represented by the graphs shown in Figure 2. For example, if XDG learns that variables \( x_2 \cdot x_3 \cdot x_4 \) all interact with \( x_1 \), it will assume that the following pairs also interact: \( (x_2, x_3), (x_3, x_4), \) and \( (x_2, x_4) \). This can have implications on decomposition of overlapping functions. XDG uses this strategy to reduce the number of function evaluations in the detection phase; however, it still requires slightly less than \( n^2 + n \) function evaluations, which is significantly more than what is needed by DG2 [58]. GDG also does not have the ability of detecting overlapping functions. Even if the algorithm is modified to check all pairs of variables, it would require \( n^2 + n \) function evaluations.

B. Comparative Analysis of Grouping Accuracy

Table I contains the summary statistics for the grouping accuracy of CCVIL, DG, GDG, XDG, and DG2 on the CEC’2010 and the CEC’2013 large-scale benchmark suite. In this paper, we use the metric proposed by Mei et al. [37], which consists of three measures: \( \rho_1 \) (interaction), \( \rho_2 \) (independence), and \( \rho_3 \) (interaction and independence). The statistics are taken over all the functions in each benchmark suite. The detailed results for individual functions can be found in Tables S-II and S-III (supplementary document, Section S-IV). The success rate indicates the number of functions for which the correct decomposition is identified. The overlapping functions are not counted since their optimal decomposition is unknown. It should be noted that XDG and CCVIL start with an interaction structure matrix of all zeros (\( \Theta = 0 \); full separability assumption). Therefore, if a pair of variables are not checked for interaction, the relevant entry of \( \Theta \) assumes its default value for the calculations of the \( \rho \)-metrics.

Table I shows that DG2 outperforms all other decomposition algorithms on the CEC’2010 and CEC’2013 suites. It is notable that the difference is more pronounced on the more difficult CEC’2013 benchmark suite. The performance of DG2 appears to be slightly lower than some variants of GDG according to \( \rho_2 \) (measure of independence). However, this is caused by three instances of the Ackley function, which affects the mean values in Table I. This behavior can be seen in Tables S-II and S-III. It should be noted that the Ackley function is not additively separable [64], which is correctly identified by DG2. However, in the benchmark suites these functions are reported as separable according to Definition 1.

Overall, DG2 shows better generalizability over a wider range of functions than all other decomposition algorithms. CCVIL shows the worst performance, while DG shows the strongest sensitivity to its control parameter (\( \epsilon \)), especially on the CEC’2010 benchmarks. This can be attributed to its static choice of \( \epsilon \). The problem with this approach is that it ignores the fact that the magnitude of the computational error in \( \lambda = |\Delta(1) - \Delta(2)| \) is correlated with the magnitude of the objective function. Therefore, on some functions, when \( \epsilon \) is smaller than the inherent computational errors, some separable variables will be considered as nonseparable. This is why increasing \( \epsilon \) generally results in a lower \( \rho_1 \) and a higher \( \rho_2 \) for DG. Unlike DG, GDG sets \( \epsilon \) proportional to the magnitude of the objective function. This is based on the rationale that a higher objective function value results in a high computational error in \( \lambda \). Equation (8) is a very simple way of choosing \( \epsilon \) proportional to the computational error:

\[
\epsilon = \alpha \cdot \min \{ f(x_1), ..., f(x_k) \},
\]

where \( x_1, ..., x_k \) are \( k \) random sample points.

Table I clearly shows that the method used by GDG is effective on the CEC’2010 benchmarks, but does not generalize well on the CEC’2013 benchmarks. To understand the reason for this behavior, the \( \epsilon \) values calculated by GDG and DG2 for selected functions are reported in Table II. Since DG2 uses a

\[\text{Also see Table S-VI of the supplementary document.}\]
different $c$ to detect interaction between each pair of variables, we report the overall mean and the median of all $c$ values as indicators. Table II clearly shows that GDG systematically overestimates the computational error which results in large $c$ values. By comparing Equations (3) and (6) with (8), we can see that the calculation of $c$ in both DG2 and GDG is a function of the objective function value. However, GDG differs in two major ways. Firstly, the constant $\alpha = \{10^{-8}, 10^{-9}, 10^{-10}\}$ is significantly larger than both $\gamma_1 \approx 2.2204 \times 10^{-16}$ and $\gamma_2 \approx 0.5108 \times 10^{-15}$ for $n = 1000$, which results in overestimation of $c$ by GDG. Secondly, $k$ sample objective function values used by GDG to detect all interactions, whereas in DG2 the quantities used in Equations (3) and (6) are only those which are involved in the calculation of $\lambda$ for a particular pair of variables. These differences contribute to the overestimation of $c$ by GDG which explains its high accuracy of detecting separable variables ($\rho_2$) at the expense of a low interaction detection accuracy ($\rho_1$). When $c$ is set to a large number, the algorithm has a tendency to classify most variables as fully separable. Conversely, a high value for $c$ makes the algorithm insensitive to weak interactions. Therefore, the algorithm may treat many weakly interacting variables as fully separable. This behavior is magnified on the CEC’2013 benchmarks due to the imbalance in the contribution of each component to the overall objective value.

### C. Optimization Results

In this section, we investigate the effectiveness of DG2 when it is used as a decomposition algorithm within a cooperative co-evolutionary framework. The empirical results are based on the CEC’2013 benchmark suite [65]. Finally, we show that in conjunction with an accurate decomposition, a contribution-based cooperative co-evolutionary algorithm shows comparable results to the state-of-the-art algorithms.

#### a) Performance comparison of decomposition methods

Table III contains the summary of the experimental results to compare the performance of DG2, DG, XDG, CCVIL, and ideal grouping, within a co-evolutionary framework. We use two different co-evolutionary frameworks: the canonical cooperative co-evolution framework in which all components are optimized in a round-robin fashion, and a contribution-based framework in which components with higher contribution to the overall solution quality are given more resources (based on the CBCC3 algorithm [66]). In this work, the $p_t$ parameter of CBCF3 is set to zero. The component optimizer of both frameworks is SaNSDE [67], and the population size of all algorithms is set to 50 as suggested by Yang et al. [67]. The maximum number of fitness evaluations is set to $3 \times 10^6$ as suggested by Li et al. [65]. Except for ideal grouping, the number of objective function evaluations used in the decomposition stage is deducted from the maximum available evaluations (a complete table is included in the supplementary document). All experimental results are based on 25 independent runs. To test the statistical significance of the results, DG and ideal grouping are compared with the baseline (DG2) using a two-tailed Wilcoxon rank-sum test with $\alpha = 0.05$.

Table III clearly shows that DG2 has an overall better performance than the other decomposition methods when it is used in a cooperative co-evolutionary framework. This is the case on both the canonical cooperative co-evolution and the contribution-based cooperative co-evolution frameworks. When comparing DG2 with ideal grouping, we can see that ideal grouping performs better on 6 functions and performs statistically similar on 9 functions, when the canonical co-operative co-evolution is used. It should be noted that the comparison between the ideal grouping and DG2 is unfair because the ideal grouping is manually given to optimization algorithm, which results in it having access to 500051 extra function evaluations. In spite of this difference, Table III shows that DG2 managed to perform as well as the ideal case on 9 functions (60% of the functions). This difference is even tighter on partially separable functions ($f_4 - f_{11}$) where DG2 performs worse than the ideal case on only one function. This clearly shows the benefit of first using some portion of the available computational resources to find an accurate decomposition of the problem before carrying out optimization. Unlike the canonical cooperative co-evolution, the difference of DG2 and the ideal grouping is wider on a contribution-based framework. Table III shows that DG2 is outperformed by the ideal grouping on 11 out of 15 functions. Since the final grouping of DG2 and the ideal grouping is identical for most of the functions, the difference can be attributed to the extra function evaluations which is available to the ideal case.

#### b) How the grouping accuracy affects the contribution-based cooperative co-evolution

To investigate the effect of an accurate decomposition on the performance of contribution-based cooperative co-evolution:
TABLE IV: Performance of contribution-based and canonical cooperative co-evolution using different decomposition methods. The number of wins, ties, and losses of the contribution-based framework against canonical cooperative co-evolution is reported. DG2 has the best improving effect on the contribution-based framework (as well as ideal grouping), especially on the partially separable functions (C2 and C3).

| Categories | CCVIL w/t | DG w/t | DG2 w/t | XDG w/t | Ideal w/t |
|------------|-----------|--------|---------|---------|-----------|
| C1: f1, f2 | 1/1/1     | 0/3/0  | 0/3/0   | 0/3/0   | 0/3/0     |
| C2: f3, f4 | 1/0/3     | 1/1/2  | 4/0/0   | 2/0/2   | 4/0/0     |
| C3: f5, f6 | 3/1/0     | 2/1/1  | 4/0/0   | 3/0/1   | 4/0/0     |
| C4: f7, f8 | 2/2/0     | 2/2/0  | 0/4/0   | 0/4/0   | 0/4/0     |
| Total      | 7/4/4     | 5/7/3  | 8/7/0   | 5/7/3   | 8/7/0     |

TABLE V: Experimental results of the contribution-based framework with DG2 (CBC3-DG2), MOS, and CMA-ES on the CEC’2013 large-scale benchmark suite using 25 independent runs. The highlighted entries are significantly better (Wilcoxon rank-sum test with Holm p-value correction, α = 0.05). DG2 allows the contribution-based framework to perform as well as the state-of-the-art even using a mediocre component optimizer.

| stats. | CBC3-DG2 | MOS [21] | CMA-ES [68] | MA-SW-Chains [22] |
|--------|----------|----------|--------------|-------------------|
| f5     | Mean     | 7.72e+06 | 1.18e+06     | 2.15e+06          |
|        | Median   | 2.11e+06 | 6.79e+06     | 1.81e+06          |
| f6     | Mean     | 2.14e+06 | 6.94e+06     | 1.87e+06          |
|        | Median   | 2.14e+06 | 6.94e+06     | 1.87e+06          |
| f7     | Mean     | 2.95e+07 | 9.03e+05     | 3.13e+05          |
|        | Median   | 2.95e+07 | 9.03e+05     | 3.13e+05          |
| f8     | Mean     | 3.39e+07 | 8.02e+07     | 2.51e+07          |
|        | Median   | 3.39e+07 | 8.02e+07     | 2.51e+07          |

A closer look at Table IV shows that DG2 achieves most of the ties on fully separable functions (C1) and overlapping and nonseparable functions (C4) for which no decomposition is done. Although DG2 can find the entire interaction structure matrix, no decomposition is performed because unlike partially separable functions, the decomposition of these functions is not unique. Therefore, both the contribution-based and the canonical cooperative co-evolution frameworks reduce to SaNSDE which is the component optimizer of both frameworks. An interesting exception is the behavior of DG and CCVIL on C4 where the contribution-based framework outperforms the canonical cooperative co-evolution framework on two cases. This is not the case for DG2, XDG, and the ideal grouping. It should be noted that DG and CCVIL decompose some functions in that category into smaller components, all of which are overlapping functions. This shows that decomposition of overlapping functions can be beneficial. This observation suggests that, by using DG2 we can learn the exact interaction pattern of the variables and identify the shared decision variables between the components in order to devise an effective decomposition for overlapping functions. However, with DG and CCVIL this is done arbitrarily depending on the order in which the variables are visited and their interaction pattern.

On the partially separable functions (C2 and C3), the contribution-based framework outperforms the canonical cooperative co-evolution when DG2 and the ideal grouping were used. This is not the case with other decomposition algorithms. Overall, the results in Table IV show that decomposition accuracy can affect the optimization performance. The experimental results suggest that the contribution-based framework requires relatively accurate decomposition in order to estimate the contribution of each component. This observation is consistent with the sensitivity analysis conducted by Kazimipour et al. [69]. In general, the contribution-based family of algorithms are sensitive to grouping noise, but in the worst case they perform as well as the canonical cooperative co-evolution which makes them a safe choice for black-box problems [69]. It should be noted that, Kazimipour et al. [69] used uniform grouping noise in their study, which equally affects both strong and weak interactions. However, we learned in Section IV-B that DG2’s grouping error is mostly attributed to detecting weakly interacting variables. This suggests that 100% accuracy is not essential in order to benefit from a divide-and-conquer scheme, but it is important on which variables does the decomposition algorithm commits the errors. An example of such a case is $f_8$ for which DG2 treated two components with weakly interacting variables as fully separable (Table S-V). However, this did not affect the overall optimization performance as reflected in Table III. Further analysis of this case is given in the supplementary document (Section S-V).
TABLE VI: Sensitivity analysis of DG2 on various complexity classes on the CEC’2013 large-scale benchmark suite. DG2 behaves similarly when growth rate of floating-point operations is not overestimated (linear and quadratic cases). However, DG2 starts to overestimate the roundoff errors when a cubic growth is assumed, which causes it to treat weakly interacting variables as separable.

| Fun. | $O(n)$ | $O(n^2)$ | $O(n^3)$ |
|------|--------|---------|---------|
| $f_1$ | 100 | 100 | 100 |
| $f_2$ | 100 | 100 | 100 |
| $f_3$ | 100 | 100 | 100 |
| $f_4$ | 99.97 | 100 | 99.99 |
| $f_5$ | 99.98 | 50.45 | 51.30 |
| $f_6$ | 100 | 100 | 100 |
| $f_7$ | 70.72 | 98.01 | 70.07 |
| $f_8$ | 99.99 | 100 | 99.99 |
| $f_9$ | 99.93 | 100 | 99.97 |
| $f_{10}$ | 99.95 | 100 | 99.95 |
| $f_{11}$ | 100 | 100 | 100 |
| $f_{12}$ | 100 | 100 | 100 |
| $f_{13}$ | 100 | 100 | 100 |
| $f_{14}$ | 99.97 | 100 | 99.99 |
| $f_{15}$ | 100 | 100 | 100 |

Success rate | 9 | 9 | 7 |

Median | 100 | 100 | 100 | 100 | 100 | 99.8 | 100 | 100 |

Mean | 97.5 | 89.3 | 90.0 | 97.4 | 89.6 | 90.2 | 95.4 | 90.6 | 91.0 |

Std. | 8.4 | 28.9 | 27.9 | 8.6 | 28.5 | 27.5 | 10.8 | 27.4 | 26.4 |

TABLE VII: Detailed grouping matrix of DG2 on $f_{10}$ based on $O(n^3)$ for estimating $\epsilon_{sup}$. The rows indicate the groups formed by DG2 and the columns represent the permutation groups from which the variables in each group were extracted.

| Groups | Size | 50 25 50 25 100 100 25 25 100 25 100 25 100 25 100 25 100 25 | 100 25 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.0 0.0 |
|--------|------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|---------------|
| P1     | 50   | 50 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P2     | 25   | 25 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P3     | 50   | 50 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P4     | 25   | 25 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P5     | 50   | 50 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P6     | 25   | 25 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P7     | 50   | 50 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P8     | 25   | 25 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P9     | 50   | 50 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P10    | 25   | 25 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P11    | 50   | 50 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P12    | 25   | 25 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P13    | 50   | 50 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P14    | 25   | 25 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P15    | 50   | 50 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P16    | 25   | 25 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P17    | 50   | 50 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P18    | 25   | 25 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P19    | 50   | 50 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| P20    | 25   | 25 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

Total | 985 | # separable variables = 1000 - 985 = 15

D. Sensitivity Analysis of DG2

It was mentioned in Section III-B that the exact calculation of the least upper bound ($\epsilon_{sup}$) is not possible due to the black-box nature of the objective function. To alleviate this problem, we proposed to estimate the number of floating-point operations based on assumptions about the complexity class of the objective function. We argued that most of the numerical calculations that do not involve complex matrix operations are of order $O(n^2)$. Additionally, error cancellations, subtraction of close numbers (Sterbenz’s Theorem [63]), and the fused multiply-add operation make the actual computational error much lower than the worst case scenario. In the previous section, we assumed a linear complexity. Here, we provide empirical results based on quadratic and cubic complexity classes to investigate the robustness of DG2 with respect to deviations from our initial assumption.

a) sensitivity to complexity of the objective function:

Table VI shows the grouping accuracy of DG2 based on different complexity classes. We can see that DG2 behaves similarly when linear and quadratic complexity classes are assumed for the number of floating-point operations. However, when a cubic complexity class is assumed, the grouping accuracy drops. From Section III-B, we know that the assumption...
increases. Our detailed analysis on P19, and P20) are the top five components with the smallest weight. The remaining 10 variables are grouped into three smaller groups. The detailed grouping matrix of P13 contains 100 variables; 55 of which are detected in G08, variables which mostly belong to P13. Table VII shows that when f13 is set to 4, the lowest weight (ρ3) of interacting variables (c3, c4) is very high and stable across all other components. It should be noted that the generated weights come from considering very weakly interacting variables as separable. It is interesting to note that P13 has the lowest weight (6.81 × 10^{-5}) among all other components in the CEC'2013 large-scale benchmark suite. The function f11 behaves in a similar way, but we do not include the details for the sake of brevity. Overall, Table VI shows that DG2 is not susceptible to moderate overestimation of ϵ_{sup}, but underestimation of ϵ_{sup} is less detrimental to its performance.

b) Sensitivity to the imbalance level: Next, we analyze the sensitivity of DG2 with respect to imbalance level among the components of the benchmark functions. The functions f4×f11 and f13×f14 have the following general form: \( \sum_{i=1}^{n} w_i \hat{f}_i(x_i) \), where \( w_i = 10^{cX(0,1)} \). The parameter c is a constant that determines the variance among the weights. In the CEC'2013 large-scale benchmark suite, c is set to 3. For our sensitivity analysis, we tested the performance of DG2 with \( c \in \{0, 1, 2, 3, 4\} \), the result of which is reported in Table VIII. The table shows that the overall grouping accuracy of DG2 is stable with various imbalance levels, except when \( c = 4 \). It is notable that the detection accuracy of separable variables (ρ2) is very high and stable across various imbalance levels. However, the detection accuracy of interacting variables (ρ1) drops when the imbalance level increases. Our detailed analysis on f4×f11 for \( c = 4 \) showed that the nonseparable components which are missed by DG2 are always among the components with the smallest weight. For example, the detailed grouping matrix of f9 with \( c = 4 \) (Table IX) shows that the missing components (P3, P4, P13, P19, and P20) are the top five components with the smallest weights. For simplicity, the log of the weights associated to each component is shown at the top of each column. Other functions have a similar behavior, but we do not include them in the analysis for the sake of brevity. Overall, this analysis shows that DG2 is not sensitive to moderate imbalance levels. When the imbalance level is very high, the inaccuracy of DG2 comes from considering very weakly interacting variables as fully separable. It should be noted that the generated weights when \( c = 4 \) are very extreme and rarely occur in real-world scenarios. Nonetheless, if this happens, treating very weakly interacting variables as fully separable is not detrimental to the optimization performance as we saw in the previous section.

V. CONCLUSION

In this paper, we proposed an improved version of the differential grouping algorithm. This new algorithm, DG2, has the following major advantages over its predecessor:
• Efficiency: lower computational cost, especially on fully separable functions;
• Accuracy: higher interaction detection accuracy;
• Robustness: lower sensitivity to computational roundoff errors;
• Applicability: the ability to detect objective functions with overlapping components, i.e., components that share decision variables. This makes it applicable to a wide array of continuous functions; and
• Practicality: no need for the user to specify a threshold parameter (ε); in other words, DG2 is parameter-free.

With respect to efficiency, we have shown mathematically that DG2 achieves the lower bound on the total number of function evaluations needed to test all pairs of variables for interaction. This effectively reduces the total number of required function evaluations by half. In addition to the improvements on efficiency, DG2 uses the information that is calculated in the process of applying the DG theorem to estimate a reliable threshold value (ε) that takes the computational error into account. The experimental results showed that DG2 significantly outperforms its predecessor on the CEC’2010 and the CEC’2013 large-scale benchmark suites.

Finally, we have shown empirically that in conjunction with DG2, the contribution-based cooperative co-evolution performs as well as the top performers of the CEC’2010 and CEC’2013 competition on large-scale optimization, as well as the well-known CMA-ES, on partially separable function. DG2 can also detect overlapping functions and can return a complete interaction structure matrix. However, due to the use of the connected components algorithm, it returns a single group containing all the decision variables. This limits the optimizer from exploiting the structural information that is found by DG2. Potential future research can focus on finding an effective decomposition for overlapping functions. Lack of a unique optimal decomposition for overlapping functions makes this a challenging task.

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SOFTWARE IMPLEMENTATION

The MATLAB/Octave and C++ implementations of the DG2 algorithm can be accessed from the following link: https://bitbucket.org/mno/differential-grouping2

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