Efficient Representation Learning Using Random Walks for Dynamic Graphs

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ABSTRACT
An important part of many machine learning workflows on graphs is vertex representation learning, i.e., learning a low-dimensional vector representation for each vertex in the graph. Recently, several powerful techniques for unsupervised representation learning have been demonstrated to give the state-of-the-art performance in downstream tasks such as vertex classification and edge prediction. These techniques rely on random walks performed on the graph in order to capture its structural properties. These structural properties are then encoded in the vector representation space.

However, most contemporary representation learning methods only apply to static graphs while real-world graphs are often dynamic and change over time. Static representation learning methods are not able to update the vector representations when the graph changes; therefore, they must re-generate the vector representations on an updated static snapshot of the graph regardless of the extent of the change in the graph. In this work, we propose computationally efficient algorithms for vertex representation learning that extend random walk based methods to dynamic graphs. The computation complexity of our algorithms depends upon the extent and rate of changes—the number of edges changed per update—and on the density of the graph. We empirically evaluate our algorithms on real world datasets for downstream machine learning tasks of multi-class and multi-label vertex classification. The results show that our algorithms can achieve competitive results to the state-of-the-art methods while being computationally efficient.

KEYWORDS
dynamic graph embedding; random walk; representation learning; information network; graph stream

1 INTRODUCTION
Graphs or networks are the natural representation of a collection of entities and the relationships between them. They are fundamental structures that have many examples in the real world, e.g., social networks, transport networks, financial transactions, and communication networks. Recently, several machine learning techniques have been proposed that use graph information to predict attributes of vertices, relationships, and the entire graph [19].

An effective approach for incorporating graph information into machine learning models is representation learning. Representation learning seeks to learn low dimensional vector representations for the vertices (vertex embeddings). The goal of the representation learning is to find a mapping of vertices to a vector representation such that distances between these vector representations meaningfully relate to similarities in the local structure of the vertices [41].

Recently, there has been significant research on representation learning and machine learning techniques that explicitly use graph information [15, 17, 18, 21, 29, 34, 38]. One common approach in many of the recent methods is unsupervised representation learning, which learns low dimensional vector representations for the vertices only based on the graph structure.

In order to capture the structure of the graphs efficiently, random walks have been shown to be scalable to large graphs [29]. In addition, random walks have been shown to be able to trade off structural equivalence (vertices that have similar local structure have similar embeddings) and homophily (vertices that belong to the same communities have similar embeddings) [17]. Random walks are combined with recent representation learning methods from language modelling to give high quality representations of vertices that can be used in downstream machine learning tasks such as vertex classification and edge prediction [17, 29]. In addition, random walk based methods have been extended to capture subgraph embeddings [6], and vertex representations in heterogeneous graphs [15]. Random walk based methods have been shown to have a fundamental link to matrix factorization [30].

However, most of the previous research on unsupervised representation learning is based on static graphs whereas most real-world graphs are dynamic —namely, they change over time. For example, in a social network, new users join the network (added vertices) and existing users add friendships (added edges). However, the static representation learning algorithms cannot measure the extent of the change in the graph. This presents a challenge for static and transductive representation learning algorithms such as DeepWalk [29] and node2vec [17] applied to large dynamic graphs, in that it becomes impractical and inefficient to re-learn vertex representations from scratch upon each change in the graph. Therefore, a method for learning representations of vertices that can explicitly utilise information about changes to the graph and update the vertex representations is desirable, in which the vertex representations can be incrementally learned as the graph changes and grows over time.

There has been recent research into how to use random-walk based representation learning for dynamic graphs. Nguyen et al. [28] used dynamic random walks on a temporal graph where at each step the next step is restricted to edges where the time is greater than at the previous step. The representations that were learned using these dynamic random walks improved predictive performance in several downstream machine learning tasks. Following on from this work, Winter et al. [14] investigated the use of time-directed walks on dynamic graphs for edge prediction and show that the past state of the graph can be used to predict the future edges. They found that node2vec [17] applied on static snapshots of the
graphs perform better in some cases than using temporally-directed
random walks.

In contrast to the previous works, in this work, we focus on
how unsupervised learning methods based on random walks can
be modified when a graph changes over time. While previous tech-
niques have used random-walks based methods on temporal graph
snapshots they have not shown how to utilise what was learned in
the previous snapshot to efficiently calculate the representations
in the next snapshot. We break this problem into two parts: firstly,
we look at how a pre-generated set of random walks generated on
one graph snapshot can be updated when the graph changes. We
show that simplistic methods to perform this update give random
walks that do not statistically represent the updated graph. We then
propose a general random walk update algorithm that produces
an updated set of walks that is statistically indistinguishable from
a set of random walks generated from scratch on the new graph.
Secondly, we investigate how to update vertex representations in-
crementally given the current set of random walks, by treating the
updates as a fine-tuning step in the DeepWalk and node2vec
algorithms [17, 29].

We demonstrate on multiple real-world datasets that our meth-
ods for updating the set of random walks and the resulting vertex
representations give comparable predictive accuracy for down-
stream tasks to that obtained by re-learning these representations
at each time snapshot of the dynamic graph while being much less
expensive computationally. We discuss the trade-offs inherent in
updating vertex representations, and how the computational cost
of updating random walks and vertex representations depends on
the number of edges that are added to the graph at each time step,
and on the density of the graph.

Our contributions are as follows:

• We propose an efficient algorithm that, given a graph struc-
ture change, produces an updated set of random walks statis-
tically indistinguishable from walks generated from scratch
on the updated graph. This update algorithm will be useful
for any task requiring a set of random walks on a graph that
is constantly changing.

• We test this algorithm by updating the skip-gram model
of [29] to work on dynamic graphs, reducing the cost of
calculating vertex representations (embeddings) by an order
of magnitude compared to computing the embeddings from
scratch.

• We empirically evaluate our algorithms with several real
world datasets for multi-class and multi-label classification
tasks.

2 UNSUPERVISED REPRESENTATION
LEARNING AND RANDOM WALKS

Given an undirected and unweighted graph, $G = \{V, E\}$, with
the set of vertices $V = \{v_1, \ldots, v_n\}$ and edges $E = \{e_1, \ldots, e_m\}$
the goal of vertex representation learning is to determine a set of
fixed length vectors, $z_k \in \mathbb{R}^N$, for each vertex $v_k$ such that similar
vertices are close in the representation space. This paper focuses
on methods that learn these representations using the conditional
probabilities of vertex pairs derived from random walks on the
graph [19].

A typical workflow for vertex representation learning on graphs
consists of the following steps: (i) update the graph, (ii) generate
random walks, (iii) learn vertex representations (embeddings), (iv)
train the downstream learning task (e.g., vertex classification). Steps
(ii) and (iii), i.e., generating random walks and learning vertex re-
presentations, are the most resource-intensive steps in the workflow.
Step (iv) involves taking the learned vertex representations and
using them as features for predictive tasks such as vertex classi-
fi cation or edge prediction. The vertex representations are learned
separately using an unsupervised objective function as this gives
multi-purpose vertex representations and also allows the use of a
small number of labelled vertices in the downstream prediction
task [29].

In this section, we describe how the DeepWalk and node2vec
algorithms [17, 29] on a static graph use random walks to generate
vertex pairs that are then used to find the vector representations $z_k$
for all vertices in the graph. In the following section, we will discuss
the issues that arise when applying these algorithms to streaming
graphs.

2.1 Random Walks

In general, a random walk can be modelled as $k$-th order Markov
chain, in which the state space is the set of graph vertices $V$ and
the future state depends on the last $k$ steps. For a $k$-th order walk
the transition probability only depends upon the previous $k$ vertices
visited by the walk [10].

A random walk of length $L$ starting from a vertex $v_{w(0)}$ con-
stitutes a sequence $W[w(0)] = \{w(0), \ldots, w(L-1)\}$, where $w(i) \in$
$\{0, 1, \ldots, n\}$ represents the vertex index at the $i$-th position in
the walk. In general, a $k$-th order random walk is generated by sampling
a vertex $w(i)$ given the $k$ previous vertices $w(i-k), \ldots, w(i-1)$
from the transition probability distribution:

$$p(w(i)|w(i-1), \ldots, w(i-k)),$$

(1)

which is non-zero only if there is an edge between vertices $v_{w(i-1)}$
and $v_{w(i)}$. To generate a $k$-th order walk, we must also sample the
$k$ initial vertices from another random sampling method in order
to calculate the transition probability (Eq. (1)). For first order walks
this amounts to sampling the vertices to start the walk from, for
higher order walks we also need another lower-order walk process
to generate the $k$ initial vertices.

To make things concrete, in the next sections we will only dis-
cuss first-order and second-order random walks that are used in
DeepWalk [29] and node2vec [17] respectively.

2.2 Vertex Pairs from Random Walks

Given a set of random walks, we extract vertex pairs that appear
close to each other in the walk. These vertex pairs represent the
structure of the graph and will be used in the next subsection to
learn the vertex representations.

Consider the case where we have a set of $K$ random walks $W_K =$
$\{W[w_j(0)], v_j \in \{1, \ldots, K\}\}$, that start from a set of initial vertices
$V_0 = \{w_0, w_2, \ldots, w_0(0)\}$ in the graph, where $W[w_j(0)] =$
$\{w_j(0), \ldots, w_j(L-1)\}$ is an independently created random walk
starting from the vertex $w_j(0)$. 
Given such a set of random walks, we sample vertex pairs from each random walk in a way analogous to words being sampled from sentences in the skip-gram model of [25]. Namely, for each random walk \( W = \{w(0), \ldots, w(L - 1)\} \) and for each vertex \( w(i) \) we take the \( p \) vertices before \( w(i) \) in the walk and create \( p \) pairs, giving the following set:

\[
S_p^-(j, l) = \{(w_j(l), w_j(l - p)), \ldots, (w_j(l), w_j(l - 1))\},
\]

the same is done for the \( p \) vertices after the vertex \( w(i) \), giving the following set of pairs:

\[
S_p^+(j, l) = \{(w_j(l), w_j(l + 1)), \ldots, (w_j(l), w_j(l + p))\}.
\]

Following the terminology of [25] we will call the first item in the pairs the target and the second item the context. Parameter \( p \) is called the context window size.

The generated target-context pairs for all random walks in set \( W_G \), together they form the corpus \( C_P(W_G) \) of vertex pairs. We define the set of all vertex pairs generated from the items before the target vertices as:

\[
C_p^- = \bigcup_{0 \leq j < K} \bigcup_{0 \leq l < L} S_{\text{min}(l, p)}^-(j, l),
\]

and all the pairs generated from items after the target vertices as:

\[
C_p^+ = \bigcup_{0 \leq j < K} \bigcup_{0 \leq l < L} S_{\text{min}(L - l - 1, p)}^+(j, l),
\]

finally, the complete vertex pair corpus is given by the union of both of these:

\[
C_p(W_G) = C_p^- \cup C_p^+,
\]

which are used to optimise the loss function as discussed in the next section.

### 2.3 The Skip-Gram Model

Given a corpus of target-context pairs \( C_p(W_G) \), the vertex representations are found by learning a vector representation for each vertex that when combined by a specified function approximates the probability of co-occurrence of the target and context vertices in the vertex pair corpus. In particular, the skip-gram model of [25] models the conditional probability of a vertex pair, \((v_t, v_c) \in V \times V\), by a log-linear function of the inner product between the vectors \( z_t \) and \( z_c \) representing the vertices, as follows:

\[
p(v_c|v_t) = \frac{\exp(z_c \cdot z_t)}{\sum_{v_k \in V} \exp(z_k \cdot z_t)}.
\]

The vertex representation vectors \( z_k \) for all vertices \( v_k \in V \) can be found by minimising the following cross-entropy loss function:

\[
-\sum_{(v_t, v_c) \in C_p(W_G)} \log \left( \frac{\exp(z_c \cdot z_t)}{\sum_{v_k \in V} \exp(z_k \cdot z_t)} \right).
\]

As the partition function in the denominator of Eq. 8 is a sum over all vertices, it is computationally intractable for all but the smallest graphs. Therefore, more efficient formulations are used to approximate this formulation, in particular Perozzi et al. [29] use a hierarchical softmax to approximate the partition function. In this work, as in more recent works [17, 18], we use negative sampling [26] to estimate Eq. 8, which is more efficient compared to hierarchical softmax.

### 2.4 Random Walk Length

Grover et al. [17] found that walks of length \( 80 \) gave the best cross-validated performance on a downstream vertex classification task when all other hyper-parameters are fixed. This meant that the number of walks was fixed and as the walk length is increased the total number of target-context pairs in the corpus \( C_p(W_G) \) also increases.

In contrast, we claim that shorter walks can give embeddings that have similar performance results on a downstream task when the number of walks is adjusted to keep the training corpus size the same. Table 1 shows similar performance of a downstream vertex classification task on the embeddings calculated from random walks of lengths \( l = 10 \) and \( l = 80 \). The performance is measured as the Macro-F1 scores for the multi-class classification on the Cora and CoCit datasets, where \( r \) is the number of walks starting from each vertex and \( l \) is the walk length. We keep the size of the corpus \( C_p(W_G) \) the same in the experiments by increasing the number of walks per vertex when we decrease the length of the walks. More information about the datasets and the experimental setups are given in Section 5.

However, we note that the training corpus \( C_p(W_G) \) will be affected by the length of the walks used to generate it. Specifically, for first-order random walks, the length of the walks will firstly change the unigram distribution of vertices in the training corpus, and secondly change the bi-gram distribution of vertex pairs through edge effects.

Firstly, the unigram distribution of the vertices appearing in the corpus will change with random walk length. In particular, as the length of the random walks becomes long, the singleton distribution of vertices will tend to the stationary distribution [11]. In comparison, short random walks will be dominated by the distribution of the initial vertices, typically a uniform distribution.

The unigram vertex distribution effectively alters the overall weighting of probabilities for each vertex in the cross-entropy loss function Eq. (8). The effect is for longer walks to give a comparatively higher weight to high-degree vertices, as they will appear more frequently in longer walks.

Secondly, edge effects due to the sampling method of Eq. (4) and Eq. (5) will bias the vertex pair corpus \( C_p(W_G) \) towards vertices that are closer together. This is because when the target vertex is closer than \( p \) to the start of the walk not all \( p \) vertex pairs can be sampled by Eq. (4). Similarly, Eq. (5) is biased to vertex pairs that are closer together at the end of the walk. As the length of the walk increases, the number of vertex pairs sampled with the full context window will increase and the effect of the edges will proportionally decrease. Therefore, the corpus \( C_p(W_G) \) will have a larger bias towards vertices that are closer together for short walks compared to long walks. For higher-order walks, there are similar edge effects due to the choice of the initial \( k - 1 \) nodes affecting the vertex pair distribution more for shorter walks.

In this section, we have described two ways that the length affects the generation of vertex pairs from random walks. We show that the effects of walk length are minimal when the size of the vertex pair corpus is controlled for. Furthermore, the changes to
the vertex pair corpus caused by different walk lengths can potentially be controlled for in other ways, for example by changing the distribution of the initial vertices of the random walks.

3 DYNAMIC GRAPHS

A dynamic graph can be represented as a series of undirected and unweighted graphs, \( \mathcal{G}^t = (\mathcal{V}^t, \mathcal{E}^t) \), where vertices \( \mathcal{V}^t = \{v_1^t, \ldots, v_m^t\} \), edges \( \mathcal{E}^t = \{e_1^t, \ldots, e_n^t\} \), and \( t \) is a discrete series of times. A dynamic graph can be considered as a set of updates taking the graph at time \( t \) and producing a modified graph at time \( t + 1 \). These updates consist of deleting or adding one or more vertices and edges. We represent the set of vertices and edges that are deleted from the graph between times \( t \) and \( t + 1 \) as \( D^t_{E,N} \) and \( D^t_{V,N} \), and the set of vertices and edges that are added to the graph between times \( t \) and \( t + 1 \) as \( A^t_{V,N} \) and \( A^t_{E,N} \) respectively. Therefore, the updated graph at time \( t + 1 \) can be given in terms of the vertices and edges as \( \mathcal{V}^{t+1} = (\mathcal{V}^t \setminus D^t_{E,N}) \cup A^t_{V,N} \) and \( \mathcal{E}^{t+1} = (\mathcal{E}^t \setminus D^t_{E,N}) \cup A^t_{E,N} \).

Our aim is given a corpus of random walks on a graph at time \( t \), \( \mathcal{W}_G(t) \), to update the random walks so that the updated set at time \( t + 1 \), \( \mathcal{W}_G(t+1) \), is statistically representative of the updated graph. Namely, the updated corpus of random walks at time \( t + 1 \) should be statistically the same as random walks drawn only from the current graph snapshot \( \mathcal{G}^{t+1} \).

Now, when updating random walks on graphs we consider the set of vertices which have changed directly as a result of the additions and deletions of vertices and edges. Specifically we denote the set of vertices contained in all added edges as \( \mathcal{V}(A^t_{E,N}) \) and the set of vertices contained in all removed edges as \( \mathcal{V}(D^t_{E,N}) \). There are several different cases to consider:

- **Deleted vertices with edges in \( D^t_{E,N} \) \( \cap \mathcal{V}(D^t_{E,N}) \)** will be removed from the graph and any random walks containing these vertices will be invalid.
- **Deleted vertices without edges in \( D^t_{E,N} \) \( \setminus \mathcal{V}(D^t_{E,N}) \)** will not be in any valid random walks and no change to the random walk corpus is needed.
- **Added vertices without edges in \( A^t_{V,N} \) \( \setminus \mathcal{V}(A^t_{E,N}) \)** will not be included in any random walks from other vertices.
- **All other added edges and deleted edges without removed vertices affect all random walks that include the vertices connected by these edges.**

To analyse the effect of graph changes on the random walks we define the following terms:

- **Affected vertices**: all vertices that are in the set of vertices contained in the set of added edges, \( \mathcal{V}(A^t_{E,N}) \) and the set of removed edges \( \mathcal{V}(D^t_{E,N}) \) but without the vertices in the deleted vertex set.

\[
\mathcal{V}^{t+1}_{\text{affected}} = \mathcal{V}(A^t_{E,N}) \cup \mathcal{V}(D^t_{E,N}) \setminus D^t_{V,N}
\]

- **Affected walks**: All random walks from the corpus of random walks \( \mathcal{W}_G(t) \) that contain at least one affected vertex.

3.1 Updating the Random Walk Corpus

Our goal is to update the random walk corpus so that it is indistinguishable from random walks generated on the updated graph. Therefore, the baseline that we will compare to is to re-generate random walks for the latest snapshot of the graph every time the graph is changed. This baseline random walk algorithm, the Static algorithm, is given in pseudocode in Algorithm 1.

The input data for the Static algorithm is the snapshot of the graph \( \mathcal{G}^{t+1} \), and the random walk parameters: number of walks \( r \) and walk length \( l \). Static takes all the vertices of the snapshot \( \mathcal{G}^{t+1} \) and initialises \( r \) random walks per vertex (Line 3), with that vertex as the initial step of those \( r \) random walks. The initialised walks are given to the function randomwalk (Line 4) that executes random walks of length \( l \) through sampling vertices from Eq. (1).

As this baseline algorithm requires a large amount of computation at each update, regardless of how small the numbers of added or deleted vertices and edges are, clearly a more efficient algorithm is needed. Such an algorithm would replace the minimum number of random walks in the current corpus of random walks \( \mathcal{W}_G(t) \) with new random walks such that the updated random walk corpus, \( \mathcal{W}_G(t+1) \), is statistically representative of the new graph \( \mathcal{G}^{t+1} \).

Firstly in the next section, we show the problem with naively updating random walks by appealing to a simple example and we introduce the Naïve Update algorithm.

3.2 Example Naïve Random Walk Update

To illustrate how a naive algorithm produces unbiased random walks we consider taking first-order random walks of length \( l \) starting from a vertex uniformly sampled from the graph \( \mathcal{G}^1 \) of Fig. 1a. The set of walks generated will be uniformly distributed over the walks shown in Table 2(a). Next, we generate the corpus of pairs \( C^1_{l} \) as described in Section 2.2. This set of adjacent pairs will be uniformly distributed over the pairs shown in Table 2(b).

The frequency of the vertex pairs \( C^1_{l} \) can be directly related to the transition probabilities for the vertex pairs given by Eq. (1) with
Table 2: (a) Random walks for graph $G^1$, (b) vertex pairs for the random walks of (a), (c) new walks for graph $G^2$, (d) vertex pairs for the random walks of (a) and (c).

| (a) | (b) | Pairs $C^1$ | Frequency | (c) | (d) | Pairs $C^1$ | Frequency |
|-----|-----|-------------|-----------|-----|-----|-------------|-----------|
| A B C | A B (A, B) | 1/4 | D B C | (A, B) | 3/18 |
| A B A | B (B, C) | 1/4 | D B A | (B, C) | 4/18 |
| B A B | B (A, A) | 1/4 | D B D | (B, A) | 4/18 |
| B C B | C (C, B) | 1/4 | (C, B) | 3/18 |
| C B C | D (D, B) | 3/18 |
| C B A | (B, D) | 1/18 |

$k = 1$, namely in the expectation:

$$E \frac{|(t, c)|}{|\mathcal{N}(v_t)|} = \frac{1}{|\mathcal{N}(v_t)|}, \quad \text{Eq.} \quad (10)$$

where $|(t, c)|$ is the number of pairs $(v_t, v_c)$ in the corpus $C^1$ and $|\mathcal{N}(v_t)|$ is the number of all pairs containing the vertex $v_t$ in the first position. For the graph $G^1$ the expected pair probabilities $p(A|B) \to 0.44$, and $p(C|B) \to 0.11$.

Now consider the graph $G^2$ shown in Fig. 1b which contains the same three vertices $A$, $B$, and $C$ as $G^1$ but adds a new vertex $D$ and an edge $B-D$. The set of affected vertices are $[B, D]$ and are shown as shaded in Fig. 1b. To update the walks to reflect the new vertex, we will naively generate walks of length 3 from the vertex $A$ with two possible choices for $B$ and $C$ because $D$ will arrive with two possible choices for $B$. Doing this we generate random walks that will be uniformly distributed over those shown in Table 2(c). However, the corpus of pairs for the set of walks consisting of the original walks of Table 2(a) and the walks of Table 2(c) do not represent the graph statistics. The vertex pairs generated from these walks will follow the distribution shown over the pairs shown in Table 2(d).

Theoretically the transition probability for vertices $A$, $C$, and $D$ given that the random walk is at vertex $B$ should all be $1/3$ as there are now three neighbours for that vertex. However, using Eq. (10) to calculate the expected empirical conditional probabilities for the vertex pairs we obtain highly biased probability estimates: $p(A|B) \to 0.44$, $p(C|B) \to 0.44$, and $p(D|B) \to 0.11$.

We see that simply adding random walks generated on the newly added vertices in the updated graph $G^2$ to the walks generated on the old graph $G^1$ gives biased conditional probability distributions and counts of vertex pairs that do not reflect the statistics of the updated graph. As the naive incremental random walk, we introduce a slightly more sophisticated algorithm, namely Naïve Update algorithm (see Algorithm 2), which initialises random walks for the affected vertices and creates random walks of length $l$. In the end, it updates the random walks $W^t$ by replacing the old walks by their corresponding re-generated walks and adding the new random walks for the new vertices (Line 4 in Algorithm 2).

However, as explained in the example, the Naïve Update algorithm does not give the same statistics of the random walks as rerunning the static algorithm Static on the current state of the graph $G^{t+1}$. In other words, the random walks given by the Naïve Update algorithm have biased empirical transition probabilities and thus do not represent the statistic structure of graph $G^{t+1}$.

In the next section, we present an algorithm that updates a set of random walks such that all statistics derived from it are identical to generating random walks on the new graph from all vertices.

### 4 DYNAMIC REPRESENTATION LEARNING ALGORITHMS

In this section, we introduce an algorithm that updates the random walks on a dynamic graph to maintain the statistical properties of the random walks as compared to sampling the random walks from scratch on the current static snapshot of the graph. We also discuss how the vertex representations learned by the skip-gram model should be updated given the vertex representations of the previous state of the graph, and the updated random walk corpus.

#### 4.1 Unbiased Random Walk Updates

In the previous section, we showed that the baseline Static algorithm, namely re-generating all walks on every graph change regardless of the extent of the change, is inefficient and incurs unnecessary computation. We additionally showed that the Naïve Update algorithm produces a random walk corpus that does not match the statistical properties of the updated graph.

To motivate an algorithm that allows an update to our random walk corpus $W_G$, we return to the definitions of affected vertices and affected walks given in Section 3 and consider a first-order random walk. In order to update a corpus of random walks we consider a random walk that has arrived at a vertex $v$, if this vertex is not affected there is no change to the neighbours of this vertex; therefore, the choice of next vertex in the random walk is unchanged. However, if the vertex $v$ is affected then the neighbours of the vertex have changed and the random walk is biased from this point as it has not considered the correct transition probabilities at the affected vertices. A corollary is that only random walks that

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**Algorithm 1 Static**

```
procedure Static($G^{t+1}$, $r$, $l$)
1: $\mathcal{V} \leftarrow \text{vertices($G^{t+1}$)}$
2: $W \leftarrow \text{initWalks($\mathcal{V}_{\text{affected}}^{t+1}$)}$
3: $W^{t+1} \leftarrow \text{randomWalk($W$, $l$, $G^{t+1}$)}$
4: return $W^{t+1}$
end procedure
```

**Algorithm 2 Naïve Update**

```
1: procedure Naïve Update($G^{t+1}$, $W^t$, $\mathcal{V}_{\text{affected}}^{t+1}$, $r$, $l$)
2: $W \leftarrow \text{initWalks($\mathcal{V}_{\text{affected}}^{t+1}$)}$
3: $W^{t+1} \leftarrow \text{randomWalk($W$, $l$, $G^{t+1}$)}$
4: $W^{t+1} \leftarrow \text{update($W^t$, $W^{t+1}$)}$
5: return $W^{t+1}$
6: end procedure
```

**Algorithm 3 Unbiased Update**

```
1: procedure Unbiased Update($G^{t+1}$, $W^t$, $\mathcal{V}_{\text{affected}}^{t+1}$, $r$, $l$)
2: $\mathcal{V}_n \leftarrow \text{newVertices($\mathcal{V}_{\text{affected}}^{t+1}$)}$
3: $\mathcal{V}_e \leftarrow \text{existingVertices($\mathcal{V}_{\text{affected}}^{t+1}$)}$
4: $W_{\text{affected}} \leftarrow \text{filter($W^t$, $\mathcal{V}_{\text{affected}}^{t+1}$)}$
5: $W_e \leftarrow \text{trim($W_{\text{affected}}$, $\mathcal{V}_e$)}$
6: $W_n \leftarrow \text{initWalks($\mathcal{V}_n$)}$
7: $W \leftarrow W_e \cup W_n$
8: $W^{t+1} \leftarrow \text{randomWalk($W$, $l$, $G^{t+1}$)}$
9: $W^{t+1} \leftarrow \text{update($W^t$, $W^{t+1}$)}$
10: return $W^{t+1}$
11: end procedure
```
contain affected vertices, namely the affected walks, need to be updated in the random walk corpus.

We propose the Unbiased Update algorithm, which is represented in Algorithm 3. The Unbiased Update algorithm updates only the affected walks by re-sampling these walks starting from the first appearance of the affected vertex. Specifically, the Unbiased Update algorithm finds the affected walks (Line 4) and trims them to the first affected vertex (Line 5). After that, Unbiased Update resumes the trimmed random walks until they are of the given length $l$.

Due to the computational expense of searching for the affected vertices in the affected walks, we also introduce the Fast Update algorithm that is the same as the Unbiased Update algorithm but instead of trimming the affected walks it re-generates the affected walks from their first step. However, as we show in the rest of this section, re-generating the affected walks from their first vertices create walks that are biased toward the walks that do not visit the affected vertices.

We have claimed that the Naïve Update and Fast Update algorithms produce random walks that have biased statistics compared to fully updating the random walks. We now show this empirically by plotting the difference between the empirical transition probabilities of the pairs generated by the four random walk update algorithms calculated using Eq. (10) and the theoretical transition probabilities on the graph $G^{t+1}$. These differences are shown in Fig. 2 which depicts the error in the transition probabilities and the normalised number of random walks (re)-generated by each algorithm compared to the number of random walks in the complete corpus $W_{G^{t+1}}$. The errors are shown in terms of mean and maximum absolute errors for all the vertices when generating random walks on the Cora dataset. The error bars are the standard deviation over 5 runs of each experiment. At each run, the graph is initialized with 10% of the edges and 5 edges are added to the graph at each time step to build the next snapshot. The details about the datasets and the experiment setup are explained in Section 5. For each snapshot of the graph, we run the four algorithms and compute the errors. We can see that the Static and Unbiased Update algorithms have a very small error compared to the Fast Update and Naïve Update algorithms which are seen to give biased transition probabilities. We note that Naïve Update has the highest error because it only generates random walks for new and affected vertices and does not consider the affected walks. We also note that the Unbiased Update algorithm can update the random walk corpus by updating less than 4% of the random walk corpus compared to Static. In the rest of the experiments, we ignore Fast Update as it has no clear advantage over Unbiased Update except that it does not require searching the random walks for the first affected vertex. However, we note that by storing the random seed that is used to generate each random walk we can regenerate the random walks up to the first affected vertex by simply reusing this seed when the walk is regenerated. Hence, by storing the random seeds for each walk we can implement the Unbiased Update algorithm with the same computational cost as the Fast Update algorithm.

4.1.1 Complexity Analysis. Consider a snapshot of graph $G^t = (V^t, E^t)$ at time $t$ with $n = |V^t|$ and $m = |E^t|$ and a random walk corpus $W_{G^t}$ consisting of $r$ random walks starting at each vertex of the graph and having length $l$. The computation cost of the Static algorithm to generate random walks for all the vertices in $G^t$ is $n \times r \times l$ steps of the random walk algorithm.

In Unbiased Update, the major functions are $\text{filter}$, $\text{randomwalk}$, and $\text{update}$ (see Algorithm 3). For simplicity, we ignore the time complexity of the $\text{trim}$ function and instead we assume all walks have the same length of $l$ for the $\text{randomwalk}$ function. The computation complexity of $\text{filter}$ using a naive linear search is $O(a \cdot n \cdot l \cdot r)$, where $a = |V^t|_{\text{affected}}$ is the number of affected vertices. However, $\text{filter}$ can be implemented faster by using advanced search techniques such as reverse-indexes on vertices. Note that the $\text{filter}$ function is embarrassingly parallel and it can be done efficiently on cluster-computing frameworks such as Apache Spark [40]. The time complexity of $\text{randomwalk}$ is $O(W \cdot l)$, where $W$ is the number of affected and new walks and $r \leq W \leq n \cdot r$. It is important to note that $l, r \ll n$. Usually, $l$ and $r$ are in the order of tens while $n$ is in the order of millions to billions, which is several orders of magnitude larger than $l$ and $r$. The smallest update on a graph is when a vertex is updated which incurs $r$ number of walks with length $l$. However, the upper bound $n \cdot l \cdot r$ in Unbiased Update, is a trivial upper bound that occurs when all the walks on the graph $G^t$ are affected, namely all vertices have an edge that has been either added or deleted. For Naïve Update, the computation cost depends only on the affected vertices, i.e., $O(a \cdot l \cdot r)$.
4.2 Updating Vertex Representations

Now we consider the problem of updating the vertex representations to give new representations that capture the current state of the graph. Specifically, we define the problem as updating the vertex representations given an updated corpus of random walks \( W_{G^{t+1}} \) and the previous representations of the vertices in the graph \( z^t_k \) for \( k \in \{0, \ldots, n(t)\} \) that were generated on the previous graph snapshot \( G^t \).

The baseline algorithm to learn vertex representations is to optimize the objective function Eq. (8) using the skip-gram model of [25]. The algorithm proceeds by initialising the vertex representations randomly. Then, the algorithm performs stochastic gradient descent to optimise the function Eq. (8) or its negative sampling version over the training data consisting of the target-context pairs in the corpus \( C_p(W_{G^{t+1}}) \) created by the updated random walks \( W_{G^{t+1}} \). For simplicity, we call this baseline algorithm Retrain.

However, Retrain is computationally expensive and inefficient for dynamic graphs, as it needs to train the skip-gram model from a random initialisation using the entire corpus \( C_p(W_{G^{t+1}}) \) regardless of the extent of the change in the graph. In Section 4.1, we see that the number of random walks that algorithms Unbiased Update, Fast Update, and Naïve Update create depends on the extent of the change on the graph.

Therefore, we propose to use this information in training the skip-gram model as well. The Incremental algorithm starts with the vertex representations from the previous graph snapshot \( z^t \) and introduces randomly initialised vectors for the new vertices in the snapshot \( A_{V^{t+1}} \). Stochastic gradient descent is then performed to minimize the skip-gram objective function using the newly generated random walks for graph \( G^{t+1} \).

Intuitively, as the optimisation process has not seen the regenerated affected walks, using the vertex pairs generated from the updated affected walks to minimise the skip-gram objective will allow updating the vertex representations with the changes that have occurred in the graph. This is, in essence, a transfer learning process [9] for the graph where the previously learned representations are fine-tuned using the vertex pairs from the updated walks, with a suitable specified learning rate. In addition, this is highly related to few-shot learning which seeks to update a learning objective that has been previously trained on a large corpus of data with a small number of examples of a new class [31]. We reserve as future work the investigation about an approach similar to the few-shot learning approach of [31] for updating the vertex representations.

As we show in Section 5, the Incremental algorithm can learn vertex representations competitive to Retrain with much less training samples.

5 EXPERIMENTS

We have implemented the random walk algorithms and the target-context pairs generator in Scala, and the skip-gram model in Python with the TensorFlow framework [5]. We run the experiments on a machine with 8 CPUs and 150GB memory.

The datasets used in the experiments are as follows:

Table 3: Datasets: number of vertices \(|V|\), number of edges \(|E|\), Labels represent the number of categories in each dataset, Density defined as \( |E|/|V|^2 \).

| Name       | \( |V| \)  | \( |E| \)  | Labels | Density          |
|------------|--------|--------|--------|-----------------|
| Cora       | 2,485  | 5,069  | 7      | \( 1.6 \times 10^{-3} \) |
| Wikipedia  | 2,357  | 11,592 | 17     | \( 4.2 \times 10^{-3} \) |
| BlogCatalog| 10,312 | 333,983| 39     | \( 6.3 \times 10^{-3} \) |
| CoCit      | 42,452 | 194,410| 15     | \( 2.0 \times 10^{-4} \) |

- Cora [3, 24]: is based on scientific papers citations. Each paper belongs only to one category among 7 categories: Neural Networks, Reinforcement Learning, Probabilistic Methods, Genetic Algorithms, Rule Learning, Theory, and Case-Based Reasoning.
- Wikipedia [3]: is the dataset of the Wikipedia web pages from 17 categories and the links between them.
- BlogCatalog [39]: is a dataset of the Wikipedia web pages from 17 categories and the links between them.
- CoCit [1, 36]: is a subgraph of the Microsoft Academic Graph, that is a network of academic papers. The dataset is a network of papers citing other papers, in which labels represent the conference that papers were published in.

Throughout the experiments, the walk lengths and the skip-gram window size are set to 10 and 8 respectively. The number of walks and the size of vector representation for each vertex are set to 80 and 128 respectively as in previous research [29]. The mini-batch size for the skip-gram model training is set to 200. The initial learning rate used is set to 0.025 for the Cora, Wikipedia, and CoCit datasets. For the BlogCatalog dataset, we observed a better performance when the learning rate is set to 0.2. We set the number of epochs to 3 for the Cora and Wikipedia datasets, but 10 and 4 for the BlogCatalog and the CoCit datasets respectively.

5.1 Vertex Classification

We evaluate the performance of the proposed dynamic graph algorithms for learning vertex representations, i.e., the Unbiased Update and Naïve Update algorithms using the Incremental method to update the skip-gram model. We compare these dynamic representation learning algorithms against our implementation of DeepWalk [29]/node2vec [17], which re-generates all the random walks and re-trains the skip-gram model. Specifically, DeepWalk and
node2vec are equivalent to the Static algorithm combined with the Retrain method.

Network snapshots in a dynamic graph often contain multiple disconnected components. Typically, for many real-world social network datasets such as the DBLP co-authorship network [2] the snapshots consist of a connected component with the majority of vertices and many smaller disconnected components. We have also observed the same behaviour when adding edges in random batches for all the datasets used in this paper. As the node representations learned on disconnected components do not have any relation to each other any downstream task must be performed separately on each disconnected component [19]. Therefore, while at every step we learn embeddings for all the components in the current snapshot, we train and evaluate the classifier only for the nodes of the largest connected component.

We evaluate the performance of the downstream classification task by using cross-validation on 9% labelled data and 10 times re-shuffling and splitting data into train/test sets. Evaluating the performance of the downstream classifier requires a considerable number of vertices for train/test splits. Therefore, we create the initial snapshots of the graphs by randomly selecting a proportion of edges. To that end, we initialised the Cora, Wikipedia, and CoCit graphs with 50%, 10%, and 4% of the edges and created consecutive snapshots of the graphs by adding more edges.

The results for the different methods are shown in Figure 3 which shows the performance of the downstream classification task as measured by the cross-validated Macro-F1 score. The performance scores are shown for the initial graph snapshot and for two snapshots chosen from the middle and the end of the experiment. We see that both dynamic update algorithms give representations that have competitive performance as compared to the full Static algorithm for each snapshot of the graph. In addition, we see that the Unbiased Update algorithm gives marginally better performance in the downstream task than the Naïve Update algorithm. This is consistent with the fact that the Naïve Update algorithm produces random walks that are biased in terms of the pair-wise transition probabilities, as discussed in Section 3.1.

Fig. 4 shows the total run time for each of the random walk and skip-gram training algorithms on the final snapshot of the graphs.
normalised to the run time of DeepWalk. The dynamic methods are, as expected, 9 times up to 160 times faster than DeepWalk (Fig. 4).

5.2 Update Rate

As we explain in Section 4.1, the computational complexity of the Naïve Update random-walk algorithm is proportional to the number of affected vertices in the graph update. This is compared to the Unbiased Update algorithm where the computational complexity is proportional to the number of affected walks. As the number of affected walks will be greater than the number of affected vertices the Naïve Update algorithm is computationally more efficient than the other algorithms (as can be seen in Fig. 4).

The run time of the Unbiased Update algorithm depends on the number of affected walks which in turn depends on the density of the graphs as well as the number of affected vertices. In our experiments, the number of affected vertices is approximately proportional to the update rate. To illustrate this for the datasets used in this paper, Fig. 5 shows the percentage of affected walks as a proportion of all the walks generated as the update rate increases. The graphs are initialised with 90% of the edges. As expected, increasing the update rate increases the number of affected walks and the denser a graph is the larger the number of affected walks for a given update rate. For example, the CoCit dataset has the lowest density of the datasets and hence we see that the proportion of affected walks is smaller for the same update rate as compared to the other datasets.

The effect of this is for a given update rate, Unbiased Update performs better for the graphs with low density. Therefore, dynamic large social graphs are naturally good targets to use our algorithm Unbiased Update as they usually have at least one order of magnitude lower density than the CoCit graph [4].

Fig. 6 depicts the run time of the Unbiased Update + Incremental method normalised to the run time of the Static + Retrain method. The results are for the last snapshot of the CoCit graph where the entire dataset has been streamed to the algorithms. We can see that when update rate is 5 edges per snapshot, the Unbiased Update + Incremental method is two orders of magnitude faster than the Static and Retrain method.

5.3 Random Walks and Bias in the Naïve Update Algorithm

As discussed in Section 4.1, Naïve Update generates random walks that are biased in terms of the pair-wise transition probabilities. To evaluate the effect of the biased random walks produced by the Naïve Update algorithm, we re-train the skip-gram model starting with a random initialisation of the weights and use all the vertex pairs generated from the random walk algorithms in the last snapshot of each graph, where each graph has 100% of the edges. We run each experiment 10 times and the mean Macro F1-score of the results for multi-class and multi-label classifications are presented in Fig. 7. The performance scores for the Unbiased Update algorithm are similar those for the Static algorithm. The random-walk corpus generated by Naïve Update is biased and this seems to give a negative effect on the downstream task performance for Cora and CoCit data sets (Fig. 7(a) and Fig. 7(b)) as seen in Fig. 7 (a) and (b). However, we don’t observe such negative effect of the random-walks generated by the Naïve Update algorithm on the Wikipedia and BlogCatalog datasets (Fig. 7(c) and Fig. 7(d)). This could be due to fact that the number of random walks that are updated by the Naïve Update algorithm is higher for graphs with higher density. In addition, the fact that the random walk corpus used to update the vertex representations is statistically biased does not necessarily affect the downstream task accuracy if this bias is small and localised. In the future, it would be of interest to investigate the differences in the representations that are found by the biased versus unbiased methods for different datasets.
Figure 7: The effect of bias in the random-walk corpus generated by the Naive Update algorithm on multi-class classification (for datasets Cora, Wikipedia, and CoCit) and multi-label classification (for dataset BlogCatalog). The results are shown in terms of Macro F1-score. The results are shown for different percentage of labelled data for training the classifier. The graphs are initialised with 10% of edges and the update rate is 5. We run our algorithm Naive Update for each snapshot of the graphs until the last snapshot, where it has 100% of edges. The vertex representations are learned by the Retrain method with the random walk samples generated by the Naive Update and Static algorithms on the final snapshot of each graph. Note that we run the same experiment for the Unbiased Update algorithm. As both the Static and Unbiased Update algorithms produce a corpus of unbiased random walks, we only present the results for the Static algorithm. As the results show, the Naive Update algorithm performs as good as the Static algorithm for the Wikipedia and BlogCatalog datasets. However, the performance of the Naive Update algorithm is slightly inferior to the Static algorithm for Cora and CoCit. Furthermore, in all the experiments the results are stable as we change the amount of training data.

### 6 RELATED WORK

Feature engineering has been long studied for graph analysis tasks. It requires domain experts to extract features [20] for vertices by handcrafting features and using feature extraction techniques [33]. In contrast, the focus of this paper is on the general-purpose representation learning approaches.

Representations learning for graphs is an important problem due to its application for tasks such as link prediction [7] and vertex classification [37]. Many of the initially proposed techniques [8, 32, 35] are to learn graph representations based on a spectral analysis of the adjacency matrix.

More recently the first and second order proximity of vertices have been used to learn vertex representations by [34] by modelling the joint probabilities and the transition probabilities of connected vertices. Driven by recent advancements in word embedding [25, 26] there have been a number of new methods to learn vertex representations in a similar way by using random walks in the graph as sentences in the same way as sentences in the word embedding methods [17, 29]. In addition, there are methods [13, 15] that extend these methods to heterogeneous graphs. A recent method VERSE [36] proposes a more flexible approach for similarity notion than of local neighbourhood. VERSE explicitly learns any similarity measure among vertices, such as personalised PageRank, to learn vertex representations.

Another approach to learning vertex representations is GraphSAGE [18] which learns a neural-network that transforms the features of a vertex and a sampled sub-graph around that vertex to a vertex representation. Specifically, GraphSAGE takes an inductive approach that can be used to generate vertex representations for vertices that are not in the training graph. Other methods offer a similar approach that requires graphs that have features for each vertex [22, 23].

We note that all of the aforementioned methods are only for static graphs. Du et al. extends LINE [34] to support dynamic graphs. Neighbourhood sampling methods based on random walks, that are the focus of this paper, have been shown to perform better than one-step and two-steps sampling in LINE [17, 29]. Chang et al. [12] propose a real-time recommender system on streaming data. In this paper, we focus on graph data and vertex representations that can be used for different downstream learning tasks.

Recently, a framework based on a modification of DeepWalk for temporal graphs has been proposed [27] that uses temporal random walks where at each step the walker is only allowed to move along an edge that has a later time than the edge used to arrive at the current vertex. The framework of [27] aims to use temporal information after the fact, whereas in this paper we focus on updating vertex representations given new edges and vertices not previously available.

### 7 CONCLUSION

Many of the real-world graphs are dynamic and change over time. However, the contemporary methods for unsupervised representation learning of vertices are mainly for static graphs. In this paper, we focused on efficient representation learning methods based on random walks for dynamic graphs. We discussed that naive incremental update of random walks results in random walks that statistically do not represent the graph structure when the graph changes over time. We proposed an intuitive way to capture changes in a dynamic graph based on the notions of affected vertices and affected walks. Following that, we proposed an incremental random walk algorithm, namely Unbiased Update, and an incremental method for representation learning, namely Incremental, which their computation cost depends on the extent of the change in the graph. Incremental random walks generated by the Unbiased Update algorithm are statistically indistinguishable from re-generating random walks by the Static algorithm. Through extensive experiments on real-world graphs, we showed that our incremental algorithms can achieve competitive results to the state-of-the-art methods while being considerably more efficient.
