Constructions of Binary Locally Repairable Codes With Multiple Recovering Sets

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ABSTRACT Locally repairable codes (LRCs) with multiple recovering sets are highly demanded in distributed storage systems. In this letter, we generalize the construction of WZL code proposed by Wang et al. and give a construction of optimal binary LRCs with multiple disjoint recovering sets which can reach the upper bound on the code rate given by Kadhe et al. Then we further generalize our idea to obtain a construction of binary LRCs with intersect recovering sets. The code rate is much higher than that of WZL code and is very close to the construction of Kruglik et al. Moreover, two special cases of this construction can reach the upper bound on the minimum distance.

INDEX TERMS Locally repairable codes, codes with availability, optimal LRCs, intersect recovering set.

I. INTRODUCTION
Distributed storage systems use redundancy to ensure data reliability, such as replication and MDS codes. Compared with traditional \([n, k]\) MDS codes, locally repairable codes (LRCs) only need to access \(r \ll k\) active nodes to recover a failure node at the cost of a small amount of storage overhead, where \(r\) is called locality. The set of these \(r\) nodes (symbols) participating in the recovery of a failure node is referred to as a recovering set of the node.

The formal definition of LRCs was first introduced by Gopalan et al. [2]. Analogous to the classical Singleton bound, they established a tradeoff between minimum distance and locality, referred to as the Singleton-type bound. A code achieving this Singleton-type bound is called optimal. After their work, other bounds were given in [3]–[6]. A tighter upper bound on the dimension \(k\) of the LRCs depending on the alphabet size was given in [7], [8]. For more studies on the bound of LRCs, one can refer to [9]–[11]. The first breakthrough construction of optimal LRCs is given in [12] by generalizing the Reed-Solomon codes. For more constructions on optimal LRCs one can refer to [13]–[17].

However, if some nodes in the recovering set are not available, we have to find an alternative set of nodes to repair the failure node. Thus, it is desirable to have multiple disjoint sets of nodes available to repair data in each node. The number \(t\) of the disjoint sets is called availability. A code is said to have locality \(r\) and availability \(t\) if every symbol has \(t\) disjoint recovering sets, denoted as \((r, t)\)-LRCs [18]. \((r, t)\)-LRCs also support parallel reading of data, which is very effective for solving the problems of degraded reading and hot data. The first upper bound on the minimum distance of the \((r, t)\)-LRC is given in [19].

\[
d \leq n - k + 2 - \left\lceil \frac{t(k - 1) + 1}{t(r - 1) + 1} \right\rceil.
\] (1)

In [20], the authors gave a bound on the code rate of \((r, t)\)-LRCs.

\[
R := \frac{k}{n} \leq \prod_{i=1}^{t} \frac{1}{1 + \frac{1}{r^i}}.
\] (2)

This bound applies to both linear codes and non-linear codes. Kadhe et al. gave a tighter bound on the rate for \((r, 3)\)-LRCs over \(\mathbb{F}_2\) in [21].

\[
R \leq 1 - \frac{3}{r + 1} + \frac{3 \log_2(2r + 4)}{(r + 1)(2r + 3)}.
\] (3)

In 2017, Kruglik et al. [22] generalized the definition of \((r, t)\)-LRCs, allowing the recovering sets of each coordinate to intersect at most \(x\) coordinates. We refer to these
codes as \((r, t, x)\)-LRCs. This feature can increase the maximum achievable code rate \([20]\) and still meet load balancing requirements. The bound on the minimum distance of \((r, t, x)\)-LRCs is given in \([1]\). Note that this bound is also valid for the case of standard \((r, t)\)-LRCs.

\[
d \leq \min_{1 \leq i \leq k-r} \frac{q^i - q^{i-1}}{d - 1} \left( n - (k - i) - \left\lfloor \frac{k - 1 - i}{r - 1} \right\rfloor \right) \tag{4}
\]

In this paper, we generalized the construction of WZL codes \([23]\), and proposed a construction of binary \((r, t)\)-LRCs. We refer to it as \textit{Construction 1}. The code rate of Construction 1 reaches the upper bound \((3)\). Noted that our codes as \((n, r, t, x)\)-LRCs. This feature can increase the maximum achievable code rate \([20]\) and still meet load balancing requirements.

We conclude in the following lemma.

\textbf{Lemma 1} \((23)\): The incidence matrix of a 1-(n, r + 1, t) design can be taken as the parity check matrix of the \((r, t)\)-LRC with length \(n\) if its blocks \(B_1, B_2, \ldots, B_b\) satisfying the following condition:

\[
|B_i \cap B_j| \leq 1, \ 1 \leq i < j \leq b. \tag{5}
\]

where \(b = \frac{n - t}{r + t}\).

\section*{C. CODES WITH AVAILABILITY AND INTERSECT RECOVERING SET}

If the recovering set of \((r, t)\)-LRC can intersect at most \(x\) positions, this code is defined as \((r, t, x)\)-LRC.

\textbf{Definition 2:} A code is said to be \((r, t, x)\)-LRC if for any coordinate \(i \in [n]\), there exist \(t\) subsets of coordinates \(R_i^1, \ldots, R_i^t \subseteq [n] \setminus \{i\}\) such that

\[
|R_i^j \cap R_i^{j'}| \leq x, \ \forall j, j' \in [t]
\]

and for all \(j \in [t]\), \(|R_i^j| \leq r\), every pair of symbols \(\alpha, \beta \in \mathbb{F}_q, \alpha \neq \beta\) we have

\[
C_{R_i^j}(i, \alpha) \cap C_{R_i^j}(i, \beta) = \emptyset.
\]

We generalize the Lemma 1 to the following result.

\textbf{Corollary 1:} The incidence matrix of a 1-(n, r + 1, t) design can be taken as the parity check matrix of the \((r, t, x)\)-LRC with length \(n\) if its blocks \(B_1, B_2, \ldots, B_b\) satisfies the following condition:

\[
|B_i \cap B_j| \leq x + 1, \ 1 \leq i < j \leq b. \tag{6}
\]

where \(b = \frac{n - t}{r + t}\).

\textbf{Proof:} If the given conditions are satisfied, then it is obvious that each row (resp. column) of the incidence matrix of such design has \(r + 1\) (resp. \(t\)) 1s, and any two different rows can intersect at most \(x + 1\) positions. The rest of the proof is similar to Lemma 1, so we omit it here. \(\square\)

\section*{III. CODE CONSTRUCTIONS}

For any positive integers \(m, a, b\), such that \(a < b, a + b \leq m\), we define a matrix over \(\mathbb{F}_q\), denoted as \(H_q(m, a, b)\), containing \([m]_q\) rows and \([m]_q\) columns. Each row of \(H_q(m, a, b)\) is associated with an \(a\)-dimensional subspace of \(\mathbb{F}_q^m\), and each column of \(H_q(m, a, b)\) is associated with a \(b\)-dimensional subspace of \(\mathbb{F}_q^m\). For \(1 \leq i \leq [m]_q, 1 \leq j \leq [m]_q\), suppose the \(i\)-th row is associated with the subspace \(W_i\) and the \(j\)-th column is associated with the subspace \(V_j\), then the \((i, j)\)-th element \(h_{ij}\) of \(H(m, a, b)\) is defined as follows:

\[
h_{ij} = \begin{cases} 
1, & \text{if } W_i \subseteq V_j \\
0, & \text{if } W_i \not\subseteq V_j.
\end{cases} \tag{7}
\]
A. A GENERAL CONSTRUCTION OF BINARY (r, t)-LRCs

When \( n = b - 1 \), the matrix \( H_q(m, a, b) \) can be regarded as a parity check matrix of an LRC with availability. We have the following theorem.

**Theorem 1:** The code \( C \) with \( H_q(m, b - 1, b) \) as a parity check matrix is a binary \([n, k, d]\) LRC with locality \( r \) and availability \( t \) where \( n = \left[ \frac{m}{b} \right]_q, k \geq \left[ \frac{b}{b} \right]_q - \left[ \frac{m}{b-1} \right]_q, \) \( r = \left[ \frac{m-b+1}{1} \right]_q - 1, t = \left[ \frac{0}{1} \right]_q \).

**Proof:** Since the matrix \( H_q(m, b - 1, b) \) contains \( \left[ \frac{m}{b} \right]_q \) rows and \( \left[ \frac{m}{b} \right]_q \) columns, we can see the code has length \( n = \left[ \frac{m}{b} \right]_q \), and dimension \( k \geq \left[ \frac{m}{b} \right]_q - \left[ \frac{m}{b-1} \right]_q \). Each row of the matrix \( H_q(m, b - 1, b) \) is associated with a \((b-1)\)-dimensional subspace \( W \) on \( \mathbb{F}_q^m \). There are \( \left[ \frac{m-b+1}{1} \right]_q \) b-dimensional subspaces containing \( W \), so each row of \( H_q(m, b - 1, b) \) has \( \left[ \frac{m-b+1}{1} \right]_q \) 1s. Each column of \( H_q(m, b - 1, b) \) is associated with a b-dimensional subspace \( V \) on \( \mathbb{F}_q^m \). Each V contains \( \left[ \frac{0}{1} \right]_q \) b-dimensional subspaces, so each column of \( H_q(m, b - 1, b) \) has \( \left[ \frac{0}{1} \right]_q \) 1s. Therefore, \( H_q(m, b - 1, b) \) is an incidence matrix of a \( \left[ \frac{m}{b} \right]_q \left[ \frac{m-b+1}{1} \right]_q \left[ \frac{0}{1} \right]_q \)-design.

Then for any i-th row associated with a b-dimensional subspace \( V_i, i \in \left[ \frac{m}{b} \right]_q \). Since \( V_i \) contains \( \left[ \frac{0}{1} \right]_q \) b-dimensional subspaces, there are \( \left[ \frac{m}{b} \right]_q \) rows which have 1 in this column. We claim that excluding the coordinate i, supports of these \( \left[ \frac{m}{b} \right]_q \) rows are pairwise disjoint. Otherwise, assume there are two rows, say the j-th row (denoted as \( h_j \), associated with the subspace \( W_j \)) and the l-th row (denoted as \( h_l \), associated with the subspace \( V_l \)), such that \( \{i, u\} \subseteq \text{supp}(h_j) \cap \text{supp}(h_l) \) for some \( u \in \left[ \frac{m}{b} \right]_q \setminus \{i\} \). It implies that \( W_j \subseteq V_l \cap V_j \) and \( W_l \subseteq V_l \cap V_j \). Then we get \( \cup \subseteq \cup \subseteq \cup \cup \). But the union of two different \((b-1)\)-dimensional subspaces is of dimension at least b, the intersection of two different \((b-1)\)-dimensional subspaces is of dimension at most b, which leads to a contradiction. Therefore, the matrix \( H_q(m, b - 1, b) \) satisfies the conditions of Lemma 1, and this completes the proof.

**Lemma 2:** The code \( C \) with \( H_q(m, b - 1, b) \) as a parity check matrix has minimum distance \( d \geq \left[ \frac{0}{1} \right]_q + 1 \).

**Proof:** It is sufficient to show that any \( \left[ \frac{0}{1} \right]_q \) columns of \( H_q(m, b - 1, b) \) are linearly independent. In the binary case, it is equivalent to show that the sum of any s columns of \( H_q(m, b - 1, b) \) is not \( 0 \) for all \( 1 \leq s \leq \left[ \frac{m}{b} \right]_q \). W.L.G. let us consider the first s columns. We denote the b-dimensional subspace corresponding to i-th column of \( H_q(m, b - 1, b) \) as \( V_i \), and (b - 1)-dimensional subspace corresponding to j-th row of \( H_q(m, b - 1, b) \) as \( W_j \). Note that \( V_i \cap V_i \) is at most a (b - 1)-dimensional subspace of \( \mathbb{F}_q^m \) for all \( i = 2, 3, \ldots, s \) and there are \( \left[ \frac{0}{1} \right] \) (b - 1)-dimensional subspace in \( V_i \). We can find a (b - 1)-dimensional subspace, say \( W_i \) such that \( W_i \neq V_i \cap V_i \) for \( i = 2, 3, \ldots, s \) (this is possible because \( s < \left[ \frac{0}{1} \right]_q \)). This implies \( W \subseteq V_i, W \not\subseteq V_i \) for all \( i = 2, 3, \ldots, s \). Hence the sum of the first s columns of \( H_q(m, b - 1, b) \) is not 0.

B. CONSTRUCTION OF BINARY (r, t)-LRCs

When \( q = 2, a = 1, b = 2 \), the parity check matrix \( H_2(m, 1, 2) \) has some special properties. We refer to this construction as **Construction 1**.

To see its properties more directly, we need to sort all the subspaces in the following manner. It is well known that every k-dimensional subspace of \( \mathbb{F}_q^n \) is the row space of a \( k \times n \) matrix of rank k, so we can use reduced row echelon form (RREF) to represent each subspace. Then take out each row of the matrix separately, assembled into a 0-1 sequence in row order, and sort the sequence in the lexical order to get the order of the subspaces.

**Remark 1:** The matrix \( H_2(m, 1, 2) \) can be viewed in the following way. Each row of \( H_2(m, 1, 2) \) is associated with a 1-dimensional subspace of \( \mathbb{F}_2^m \), which only contains one vector of length m, so it can be seen as a binary number of length m. Similarly, each column of \( H_2(m, 1, 2) \) is associated with a 2-dimensional subspace of \( \mathbb{F}_2^m \). Since a 2-dimensional vector space contains three vectors, each vector can be seen as a binary number of length m, where the first two vectors are the first and second row of the \( 2 \times n \) RREF matrix. The third vector is the sum of the first two vectors (under \( \mathbb{F}_2 \)). If a binary number corresponding to the ith row is contained in the third binary numbers of the jth column, the value at the ith row and jth column of the matrix \( H_2(m, 1, 2) \) is 1, otherwise 0.

Below is an example of \( H_2(m, 1, 2) \).

**Example 1:** Suppose \( m = 3 \), the code \( C \) with \( H_2(3, 1, 2) \) (Fig.1) as a parity check matrix is a \((2, 3)\)-LRC with length \( n = 7 \), dimension \( k = 3 \), minimum distance \( d = 4 \). We label all vectors in the vector space at the beginning of rows and columns (regard a vector as a binary number and convert it to a decimal number).

| \( H_2(3, 1, 2) = \) | 2 | 4 | 4 | 4 | 5 | 6 |
|-------------------|---|---|---|---|---|---|
|                   | 1 | 1 | 2 | 3 | 2 | 3 | 1 |
|                   | 3 | 5 | 6 | 7 | 7 | 6 | 7 |

**FIGURE 1.** The matrix \( H_2(3, 1, 2) \).

The above example has the same parameters as the Simplex code with \( m = 3 \). Next we prove some properties of \( H_2(m, 1, 2) \) to help understand the code \( C \).

**Lemma 3:** For any positive integer \( m \geq 3 \), the matrix \( H_2(m, 1, 2) \) is of the block form

\[
H_2(m, 1, 2) = \begin{pmatrix}
H_2(m-1, 1, 2) & * \\
0 & A
\end{pmatrix}
\]  \( (8) \)
where

\[ A = \begin{pmatrix} 1 & \cdots & 1 \\ I_{m-1} & 1 \end{pmatrix} \]

\( I_{m-1} \) is the identity matrix of size \([m-1]_2 \times 0\) and \( A \) is a zero matrix. \( \ast \) represents the arbitrary matrix over \( \mathbb{F}_2 \). Particularly, \( H_2(1, 2, 2) = (1, 1, 1)^t \).

**Proof:** It is easy to see \( H_2(2, 1, 2) = (1, 1, 1)^t \). For the case \( m \geq 3 \), noticed that the RREF matrices corresponding to the former \([m-1]_2 \times 2 \) rows of the matrix \( H(m, 1, 2) \) all have 0s as their first entries, thus these rows can be regarded as 1-dimensional subspaces of \((m-1)\)-dimensional space. Since the subspaces associated with columns are sorted, the subspaces correspond to the RREF matrices whose first column is all 0s must be at the top, and these columns can be regarded as 2-dimensional subspaces of \((m-1)\)-dimensional space. There are a total of \([m-1]_2 \times 2 \) such subspaces. Therefore, the upper left block of \( H(m, 1, 2) \) (i.e. the former \([m-1]_2 \times 2 \) rows and the former \([m-1]_2 \times 2 \) columns) is the matrix \( H(m, 1, 2) \). The RREF matrices corresponding to the latter \(2m-1\) rows of the matrix \( H(m, 1, 2) \) all have 1s as their first entries, but the RREF matrices corresponding to the former \([m-1]_2 \times 2 \) columns of the matrix \( H(m, 1, 2) \) all have 0s as their first entries. As a result, the bottom left block is a \(2m-1 \times [m-1]_2 \) zero submatrix.

Regarding the matrix \( A \), note that the RREF matrix of the first row of \( A \) are \((1, 0, \cdots, 0)\), which has \( m - 1 \) zeros. Since the subspaces are sorted, the subspaces correspond to the RREF matrices whose first row is \((1, 0, \cdots, 0)\) must be ranked after the RREF matrices whose first column is 0, before other subspaces. There are a total of \([m-1]_2 \times 2 \) such subspaces. Therefore, the former \([m-1]_2 \times 2 \) columns of the first row of matrix \( A \) are all 1s. Moreover, if we add the two rows of the RREF matrix corresponding to the former \([m-1]_2 \times 2 \) columns of matrix \( A \), we can get the vectors whose binary representation from \((1, 0, \cdots, 1)\) to \((1, 1, \cdots, 1)\) (decimal representation from \(2^m-1\) to \(2^m-1\)). The vectors corresponding to the rows start from the second row of matrix \( A \) whose binary representation also from \((1, 0, \cdots, 1)\) to \((1, 1, \cdots, 1)\). Therefore, the bottom left block of matrix \( A \) is an identity matrix of size \([m-1]_2 \).

According to the block form of the matrix \( H(m, 1, 2) \), it is easy to see \( \text{rank}(H(m, 1, 2)) \geq \sum_{i=1}^{m-1} [i]_2 = 2^m - 1 - m \). In fact, the rows of \( H(m, 1, 2) \) are linearly dependent, so some rows can be deleted. We define the following set.

\[ E_i = \left\{ (2j - 1)2^i + k : 1 \leq j \leq 2^{m-i-1}, 0 \leq k \leq 2^j - 1 \right\} \]  \( \text{(9)} \)

From the definition of the set \( E_i \), if we fix \( m = 4 \), we can get,

\[ E_0 = \{1, 3, 5, \cdots, 15\}, \]

\[ E_1 = \{2, 3, 6, 7, 10, 11, 14, 15\}, \]

\[ E_2 = \{4, 5, 6, 7, 12, 13, 14, 15\}, \]

\[ E_3 = \{8, 9, 10, 11, 12, 13, 14, 15\}. \]

If all the elements in the \( E_i \) are converted into binary form, we will find that the first digit of all elements in \( E_0 \) is 1, that is, all odd numbers; the second digit of all the elements in \( E_1 \) is 1; the third digit of all elements in \( E_2 \) is 1; the fourth digit of all elements in \( E_3 \) is 1. In fact, \( E_i \) is the set of elements in the binary form of all elements in the complete set \( U = \{1, 2, \cdots, 2^m-1\} \) whose \((i+1)\)th digit is 1.

For convenience, we also define

\[ R_i = \{ j^{\text{th} \text{ row of } H(m, 1, 2) : j \in E_i} \} \]

**Lemma 4:** The rank of matrix \( H(m, 1, 2) \) is \( 2^m - 1 - m \).

**Proof:** In the following, we show that each \(2^i\)-th row in matrix \( H(m, 1, 2) \) is an \( F_2 \)-linear combination of all rows in \( R_i \) for \( i = 0, 1, \cdots, m-1 \).

Theorem 1 has showed that every column of matrix \( H(m, 1, 2) \) has three 1s. It is sufficient to show that the number of 1 in any column of the rows in \( R_i \) is 0 or 2. Since the subspaces are sorted in lexical order, the binary representation of the vector corresponding to each row can be regarded as row number, and the \((i+1)\)-th digits of the binary representation of all elements in the set \( E_i \) are all 1. For any columns in \( R_i \), say \( j^{\text{th}} \text{ column}, j \in \{1, 2, \cdots, [m/2]_2\} \). Suppose the number of 1 in \( j^{\text{th}} \text{ column} \) is 1 (resp. 3), this means there is only one (resp. three) vector whose \((i+1)\)-th coordinate is 1 in the space corresponding to the \( j^{\text{th}} \) column, which is impossible. Because the \((i+1)\)-th coordinate of the remaining two vectors in the space is 0 (resp. 1), so the \((i+1)\)-th coordinate of the third vector obtained by adding these two vectors must also be 0, which leads to a contradiction, from which the result follows.

**Lemma 5:** When \( m \geq 3 \), the code \( C \) which has the parity check matrix \( H_2(m, 1, 2) \) has minimum distance \( d = 4 \).

**Proof:** From Lemma 3, the upper left block of the matrix \( H_2(m, 1, 2) \) contains \( H_2(3, 1, 2) \) for \( m \geq 3 \), and the bottom left block is a zero matrix. The matrix \( H_2(3, 1, 2) \) has 4 columns that are linearly dependant (see Fig. 1). So there are 4 columns in the matrix \( H_2(m, 1, 2) \) that are linearly dependant. Combine with the Lemma 2, the result follows.

**Corollary 2:** The code \( C \) with \( H_2(m, 1, 2) \) as a parity check matrix is a binary \( [n, k, d] \)-LRC with locality \( r \) and availability \( t \) where \( n = [m/2]_2, k = [m]_2 - [m/2]_2 + m, r = 2^m - 2^t - 2, t = 3, d = t + 1 = 4 \). Therefore, the code is an optimal LRC with availability which reaches the bound (3).

### C. CONSTRUCTIONS OF \((r, t, x)\)-LRCs

When \( q = 2, a = 1, b \geq 3 \), we can get the LRC with availability in which the recovering sets can intersect in a small number of coordinates. We refer to this construction as Construction 2.

Let us give an example of \( H_2(m, 1, b) \).
Example 2: Suppose $m = 4$, the code $C$ which has the parity check matrix $H_2(4, 1, 3)$ is a (6, 7, 2)-LRC with $n = 15$, $k = 10$, $d = 4$.

\[
\begin{pmatrix}
1 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 \\
1 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 \\
1 & 1 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 1 \\
1 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 & 1 & 0 \\
1 & 0 & 1 & 0 & 0 & 1 & 0 & 1 & 1 & 0 & 0 & 1 \\
0 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 0
\end{pmatrix}
\]

Note that the minimum distance of the above code has reached the bound (4).

Theorem 2: The code $C$ which has the parity check matrix $H_2(1, m, 1, 1)$ is a $(r, t, x)$-LRC for $r \geq 3$, where $r = \left(\frac{m}{2} - 1\right) - 1, t = \left(\frac{m}{2} - 2\right) - 1$.

Proof: Each row of matrix $H_2(m, 1, b)$ is associated with a 1-dimensional subspace $W$ on $\mathbb{F}_2$. There are $\left(\frac{m}{2} - 1\right)$, $b$-dimensional subspaces containing $W$, so each row of $H_2(m, 1, b)$ has $\left(\frac{m}{2} - 1\right)$ 1s. Each column of $H_2(m, 1, b)$ is associated with a $b$-dimensional subspace $V$ on $\mathbb{F}_2$. Each $V$ contains $\left(\frac{m}{2}\right)$, 1-dimensional subspace, so each column of $H_2(m, 1, b)$ has $\left(\frac{m}{2}\right)$ 1s. Therefore, $H_2(m, 1, b)$ is an incidence matrix of a $1 - \left(\begin{smallmatrix}
\frac{m}{2} - 1 \\
2
\end{smallmatrix}\right) > \left(\begin{smallmatrix}
\frac{m}{2} - 1 \\
2
\end{smallmatrix}\right)$-design.

For any two rows of $H_2(m, 1, b)$, say $i$th row and $j$th row, which corresponds to the subspace $W_i$ and $W_j$ respectively, for $1 \leq i < j \leq \left(\begin{smallmatrix}
\frac{m}{2} \\
2
\end{smallmatrix}\right)$, $W_i \cap W_j$ is a 2-dimensional subspace. For any column of $H_2(m, 1, b)$, there are $\left(\frac{m}{2} - 1\right)$, $b$-dimensional subspaces which contain such 2-dimensional subspace $W_i \cap W_j$, that is the supports of any two rows of $H_2(m, 1, b)$ intersect at $\left(\frac{m}{2} - 2\right)$ coordinates, i.e. $x = \left(\frac{m}{2} - 2\right)$ 1s. Therefore, the matrix $H_2(m, 1, b)$ satisfies the conditions of Corollary 1, and it complete the proof. □

Then we give a brief analysis of the structure of the matrix $H_2(m, 1, b)$ for $b \geq 3$.

Theorem 3: For any positive integer $3 \leq b < m$, the matrix $H_2(m, 1, b)$ is of the block form

\[
H_2(m, 1, b) = \begin{pmatrix}
H_2(m - 1, 1, b) & * \\
0 & * 
\end{pmatrix},
\]

where $0$ is a zero matrix, * represent arbitrary matrix over $\mathbb{F}_2$. In particular $H_2(b, 1, 1) = (1, \ldots, 1)^T$ is a column vector of length $\left(\frac{m}{2}\right)$.

Proof: The proof is similar to the proof of Theorem 3, so we omit it here. □

Next, we generalize the method of proving the rank of $H_2(m, 1, 2)$ to prove the rank of $H_2(m, 1, b \geq 3)$. Since all the subspaces are defined over $\mathbb{F}_2$, for the convenience of proof, we also view the vector in the subspace as a binary number.

Lemma 6: For any $m$-dimensional vector space $V$ over $\mathbb{F}_2$. Let $S$ be a set of column indices such that $1 \leq |S| \leq m - 1$. Then $|\{v \in V : v_1 = 1, \forall i \in S\}| \equiv 0 \mod 2$.

Proof: Let $n \geq m$ be the vector length of $V$. We represent $V$ as a RREF matrix $G$ of size $m \times n$. Denote $G_S$ as a submatrix formed by the columns of indices in $S$. We transform $G$ to a matrix $G'$, such that $G_S$ is in its RREF. If $k$ is the rank of $G_S, 0 \leq k \leq |S|$, then the last $(m-k)$ rows of $G'$ are all zero vectors. Notice that the vector space generated by the first $k$ rows of $G'$ contains at most one all 1s vector. Now, consider the vector space generated by the last $(m-k)$ rows of $G'$, its cardinality is $2^{m-k}$, and all vectors have 0 entries at the columns of indices in $S$. Therefore, the vector space generated by $G$ contains either 0 or $2^{m-k}$ vectors that have all 1s in the columns of indices in $S$. From which the result follows. □

In fact, rows of the matrix $H_2(m, 1, b)$ are linearly dependent, so some rows can be deleted. We define the following set consisting of $s$-tuples of positive integers:

\[
C_s^m = \{ (\alpha_1, \cdots, \alpha_s) : 1 \leq \alpha_1 < \cdots < \alpha_s \leq m \}
\]

Actually, $C_s^m$ is a set of $s$-combination of $[m], |C_s^m| = \binom{m}{s}$.

Let $E_j^i$ be a set of the all $m$-bit binary numbers whose $i$-th bit is 1, for all $i \in (\alpha_1, \alpha_2, \cdots, \alpha_s)$, where $(\alpha_1, \alpha_2, \cdots, \alpha_s)$ is the $j$-th tuple in $C_s^m$. Let $(E_j^i)_{min}$ be the smallest element in $E_j^i$. Obviously, when $s$ traverses from 1 to $b - 1$, all these $(E_j^i)_{min}$ are different and

\[
\sum_{i=1}^{b-1} |E_j^i| = \sum_{i=1}^{b-1} \binom{m}{i}.
\]

For convenience, we also define

\[
R_j^i = \{ i - \text{th row of } H_2(m, 1, b) : i \in E_j^i \}
\]

Lemma 7: For any $b \geq 3$, $\text{rank}(H_2(m, 1, b)) \leq \sum_{i=1}^{b-1} \binom{m}{i}$. Therefore the code $C$ which has the parity check matrix $H_2(m, 1, b)$ has dimension $k \geq \left[ \frac{m}{2} - 1 \right] + \sum_{i=1}^{b-1} \binom{m}{i}$.

Proof: In the following, we show that each $(E_j^i)_{min}$-th row in matrix $H_2(m, 1, b)$ is an $\mathbb{F}_2$-combination of all rows in $R_j^i$ for $j = 1, 2, \cdots, \binom{m}{s}$ and $s = 0, 1, \cdots, b - 1$.

The matrix $H_2(m, 1, b)$ is defined over $\mathbb{F}_2$, it is sufficient to show that the number of 1 in any column of the rows in $R_j^i$ is even. Since the subspaces are sorted in lexical order, the binary representation of the vector corresponding to each row can be regarded as a row number. According to Lemma 6 and the definition of $E_j^i$, for each row in $R_j^i$, its $i$-th column is 1 for all $i \in (\alpha_1, \alpha_2, \cdots, \alpha_s)$, where $(\alpha_1, \alpha_2, \cdots, \alpha_s)$ is the $j$-th tuple in $C_s^m$. Therefore, the number of 1 in any column of the rows in $R_j^i$ is even. Moreover, because all these $(E_j^i)_{min}$ are different, we can sort $(E_j^i)_{min}$ in increasing order, then all $(E_j^i)_{min}$-th rows of the $H_2(m, 1, b)$ can be deleted in this order. From which the result follows. □

Corollary 3: The code $C$ which has the parity check matrix $H_2(m, 1, b)$ has dimension

\[
k \geq \left[ \frac{m}{2} \right] - \left[ \frac{m}{2} \right] + \sum_{i=1}^{b-1} \binom{m}{i}.
\]
so the code rate is,
\[ R \geq 1 - \left[ \begin{array}{c} m \\ 2 \end{array} \right] - \sum_{i=1}^{b-1} \left( \begin{array}{c} m \\ i \end{array} \right). \] (13)

Remark 2: In fact, if we use the method similar to the analysis of matrix \( H_2(m, 1, 2) \) to deeply analyze the block structure of \( H_2(m, 1, b) \), we can get that the rank of the matrix \( H_2(m, 1, b) \) is indeed \( \sum_{i=1}^{b-1} \left( \begin{array}{c} m \\ i \end{array} \right) \).

Lemma 8: The code \( C \) which has the parity check matrix \( H_2(m, 1, b) \) has minimum distance \( d \geq 4 \).

Proof: It is obvious that any two columns of the matrix cannot be equal. It is sufficient to show that the sum of any 3 columns of \( H_2(m, 1, b) \) is not 0. We denote the \( b \)-dimensional subspace corresponding to the \( i \)-th column of \( H_2(m, 1, b) \) as \( V_i \). Note that any two different subspaces \( V_i \cap V_j \) is at most a \((b - 1)\)-dimensional subspace of \( F_2^m \). And a \((b - 1)\)-dimensional subspace contains \( \left[ \begin{array}{c} b-1 \\ 1 \end{array} \right] \), so the supports of any two columns of \( H_2(m, 1, b) \) can intersect at most \( 2^{b-1} - 1 \) coordinates. However, the number of 1 in any column of \( H_2(m, 1, b) \) is
\[ \left[ \begin{array}{c} b \\ 1 \end{array} \right] = 2^b - 1 > 2 \times (2^{b-1} - 1) = 2 \times \left[ \begin{array}{c} b-1 \\ 1 \end{array} \right]. \]
This means that the sum of any 3 columns cannot be 0, from which the result follows. \( \square \)

IV. COMPARISON WITH OTHER CONSTRUCTIONS

A. GENERAL CONSTRUCTION

Our general construction is a binary regular LDPC code with girth \( > 4 \). Hao et al.’s [25] proposed a construction of LRC codes with information symbols by combining an existing regular LDPC and an identity matrix. But we directly construct the parity-check matrix to obtain an LRC code, and this matrix can be viewed as a girth matrix of BIBD with \( \lambda = 1 \), so our construction can also be regarded as a kind of BIBD-LDPC codes.

B. CONSTRUCTION 1

Among \((r, t)\)-LRCs that have the same availability \( t = 3 \) as our Construction 1, WZL code is the one that has good parameters. It has been shown that WZL code has a higher rate than that of direct product code and Prakash et al.’s construction [3].

Recall that the parameters of a WZL code are \( n_w = \left( \begin{array}{c} r+t \\ t \end{array} \right) \), \( d_w = t + 1 \), \( R_w = \frac{r}{r + t + (r + 1)(2r + 3)} \). For all \( r > 0 \), the code length of Construction 1 is \( n_1 = \frac{r}{r + t + (r + 1)(2r + 3)} \), which is shorter than that of WZL code \( n_w = \left( \begin{array}{c} r + 3 \\ 3 \end{array} \right) \). Both constructions have the same minimum distance. To compare the code rate, Construction 1 is rate optimal. Indeed, it is easy to see that we always have a greater code rate than \( R_w \) for \( r > 0 \).

C. CONSTRUCTION 2

There are few works on constructions of \((r, t, x)\)-LRCs. Kruglik et al.’s construction is based on WZL code, its parameters are \( n_k = (x + 1) \left( \begin{array}{c} t+r \\ t \end{array} \right) \), \( d_k = 2 \), \( R_k = (r + 1)(x + 1) - 1 \), \( R_k = \frac{r+(r-1)x}{r+t+(r-1)x} \), so our Construction 2 has a much shorter code length and \( 2 \times \) greater minimum distance than Kruglik’s construction. We also compare our code rate to that of Kruglik et al. and WZL code with the same locality \( r \) and availability \( t \), as shown in Fig. 2 and Fig. 3.

The figures show our code rate are greater than WZL code, but slightly less than Kruglik et al. Moreover, when \( m = 4, b = 3 \) and \( m = 5, b = 4 \), our construction reaches the minimum distance bound (4) of \((r, t, x)\)-LRC codes. Note that since the matrix \( H_2(m, 1, b) \) has the above-mentioned block form (see Theorem 3), when \( b \) is fixed, the minimum distance is also fixed, which is the same as the code has parity check matrix \( H_2(b+1, 1, b) \).

V. CONCLUSION

In this paper, we generalize the construction of WZL codes and propose two constructions of LRC codes. Construction 1 can produce optimal \((r, t)\)-LRCs, which can reach the upper bound of code rate (3). Construction 2 has much higher rate than that of WZL codes and attain the upper bound on minimum distance (4) in two special cases. Moreover, we give a sufficient condition for a 1-design’s incidence matrix that can be the parity-check matrix of a \((r, t, x)\)-LRCs.
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