A BRIEF HISTORY OF PERSISTENCE

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ABSTRACT. Persistent homology is currently one of the more widely known tools from computational topology and topological data analysis. We present in this note a brief survey on the evolution of the subject. The goal is to highlight the main ideas, starting from the subject’s computational inception more than 20 years ago, to the more modern categorical and representation-theoretic point of view.

1. The Early Years: Computer Science Meets Geometry, Algebra and Topology

Our current view of persistent homology can be traced back to work of Patrizio Frosini (1992) on size functions [29], and of Vanessa Robins (1999) [44] on using experimental data to infer the topology of attractors in dynamical systems. Both approaches rely on singular homology as a shape descriptor, which leads to what is known today as the “homology inference problem”: Given a finite set $X$ (the data) sampled from/around a topological space $X$ (e.g., the attractor), how can one infer the homology of $X$ from $X$ with high confidence? See for instance [41] for the case when $X$ is a compact Riemannian submanifold of Euclidean space, and $X \subset X$ is sampled according to the intrinsic uniform distribution. From here on out it will be useful to think of $X$ and $X$ as subspaces of a bounded metric space $(M, \rho)$. In this case, one can formalize the statement “$X$ approximates $X$” by saying that if $Z \subset M$, $\epsilon \geq 0$, and $Z(\epsilon) := \{ x \in M : \rho(x, Z) \leq \epsilon \}$, then the Hausdorff distance

$$d_H(X, X) := \inf \left\{ \epsilon > 0 : X \subset X(\epsilon) \text{ and } X \subset X(\epsilon) \right\}$$

is small. The goal is then to approximate the topology of $X$ by that of $X(\epsilon)$. Below in Figure 1 we illustrate the evolution of $X(\epsilon)$ as $\epsilon$ increases.

![Figure 1. Some examples of $X(\epsilon)$ for $X \subset \mathbb{R}^2$ sampled around the unit circle, and $\epsilon$ values $0 < \epsilon_1 < \epsilon_2 < \epsilon_3$.](image)

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In order to capture the multiscale nature of $X = \{X(\epsilon)\}_\epsilon$, and deal with the instability of topological features in $X(\epsilon)$ as $\epsilon$ changes, Frosini and Robins introduced (independently) the idea of homological persistence: for $\epsilon, \delta \geq 0$ let

$$\iota_{\epsilon,\delta} : X(\epsilon) \hookrightarrow X(\epsilon+\delta)$$

be the inclusion, and consider the induced linear map in homology with coefficients in a field $\mathbb{F}$

$$\iota_{\epsilon,\delta}^* : H_n \left( X(\epsilon); \mathbb{F} \right) \longrightarrow H_n \left( X(\epsilon+\delta); \mathbb{F} \right)$$

The image of $\iota_{\epsilon,\delta}^*$ is the $\delta$-persistent $n$-th homology group of the filtered space $X$ at $\epsilon$, denoted $H_{n,\delta}^\epsilon(X; \mathbb{F})$; and $\text{rank} \left( \iota_{\epsilon,\delta}^* \right)$ is the persistent Betti number $\beta_{n,\delta}^\epsilon(X; \mathbb{F})$.

The design of algorithms to efficiently compute/approximate these integers is of course predicated on first replacing the spaces $X(\epsilon)$ by finite, combinatorial models of their topology. Fortunately there is a vast literature on how to do this. Take for instance the Vietoris-Rips complex, first introduced by Leopold Vietoris in the nineteen-twenties in an attempt to define a homology theory for general metric spaces [47]. It is defined, for $Z \subset \mathbb{M}$ and $\epsilon \geq 0$, as the abstract simplicial complex

$$R_\epsilon(Z) := \{ \{z_0, \ldots, z_k\} \subset Z : \rho(z_i, z_j) \leq \epsilon \text{ for all } 0 \leq i, j \leq k \}$$

Below in Figure 2 we show an example of how $R_\epsilon(Z)$ evolves as $\epsilon$ increases, for $Z \subset \mathbb{R}^2$ sampled around the unit circle, and for $\epsilon$ values $0 < \epsilon_1 < \epsilon_2 < \epsilon_3$.

![Figure 2. Some examples of the Rips complex, for points sampled around the unit circle in $\mathbb{R}^2$.](image)

Notice that $R_\epsilon(Z) \subset R_{\epsilon+\delta}(Z)$ whenever $\delta \geq 0$; in other words, $\mathcal{R}(Z) = \{R_\epsilon(Z)\}_\epsilon$ is a filtered simplicial complex. Janko Latschev shows in [35] that when $X$ is a closed Riemannian manifold, there is an $\epsilon_0 > 0$, so that if $0 < \epsilon \leq \epsilon_0$, then there exists $\delta > 0$ for which $d_H(X, X) < \delta$ implies that the geometric realization of $R_\epsilon(X)$ is homotopy equivalent to $X$. Discarding the manifold hypothesis — which is not expected to hold in general applications — highlights the value of persistence as a homology inference tool. Indeed, in [17] Chazal, Oudot and Yan show that if $X \subset \mathbb{R}^d$ is compact with positive weak feature size\(^1\) [16], and $X \subset \mathbb{R}^d$ is finite with $d_H(X, X)$ small, then there exists a range for $\epsilon > 0$ where $H_{n,3\epsilon}^\epsilon(\mathcal{R}(X); \mathbb{F})$ is isomorphic to $H_n(X(\epsilon'); \mathbb{F})$. It is worth noting that while these theorems deal with small $\epsilon$, far less is known about the large-scale regime. Indeed, aside from trivial examples, the circle is (essentially) the only space $Z$ for which the homotopy type of $R_\epsilon(Z)$ is known explicitly for all $\epsilon > 0$ [1, 2].

The efficient computation of the persistent Betti numbers of a finite filtered simplicial complex $\mathcal{K} = \{K_0 \subset K_1 \subset \cdots \subset K_J = K\}$, was addressed by Edelsbrunner, Letscher and Zomorodian in (2000) [27], for subcomplexes of a triangulated 3-sphere and homology with coefficients in $\mathbb{F}_2 = \{-1, 1\}$. This restriction was a

\(^1\)this is a notion of how complex the embedding of $X$ into Euclidean space is.
tradeoff between generality and speed: the algorithm was based on previous work of Delfinado and Edelsbrunner [23] to compute (standard) Betti numbers incrementally in time $O(N\alpha(N))$, where $N$ is the number of simplices of $K$ and $\alpha$ is the inverse of the Ackermann function [19]. Since the Ackermann function grows very rapidly, its inverse $\alpha$ grows very slowly. Though limited in generality, the approach by Delfinado and Edelsbrunner highlights the following idea: If $K_j$ is obtained from $K_{j-1}$ by adding a single simplex $\tau \in K$, and $H_n(K_{j-1}; \mathbb{F}) \rightarrow H_n(K_j; \mathbb{F})$ is not surjective, then either $\tau$ is an $n$-simplex creating a new homology class, or it is an $n+1$-simplex eliminating a class from $K_{j-1}$. Thus, simplices in $K$ that either create or annihilate a given persistent homology class can be put in pairs $(\tau, \sigma)$ of the form creator-annihilator. These pairings are in fact a byproduct of the incremental algorithm of Delfinado and Edelsbrunner. The barcode is also introduced in [27] as a visualization tool for persistence: each pair $(\tau, \sigma)$ yields an interval $[j, \ell)$, where $j$ (birth time) is the smallest index so that $\tau \in K_j$, and $\ell > j$ (death time) is the smallest index for which $\sigma \in K_\ell$. Thus, long intervals indicate stable homological features throughout $K$, while short ones reflect topological noise. The resulting multiset of intervals (as repetitions are allowed) is called a barcode. The notation is $\text{bcd}_n(K)$. Moreover, the barcode subsumes the persistent Betti numbers, since $\beta_n^\delta(K; \mathbb{F})$ is the number of intervals $[j, \ell) \in \text{bcd}_n(K)$ with $j \leq \epsilon$ and $\ell > \epsilon + \delta$. Below in Figure 3 we show an example of a filtered simplicial complex, the simplicial pairings $(\tau, \sigma)$, and the resulting barcodes.

2. Here Comes the Algebra

The developments up to this point can be thought of as the computational and geometric era of persistent homology. Around 2005 the focus started to shift towards algebra. Zomorodian and Carlsson introduced in [50] the persistent homology

$$PH_n(K; \mathbb{F}) := \bigoplus_{j \in \mathbb{Z}} H_n(K_j; \mathbb{F}), \quad K = \{K_j\}_{j \in \mathbb{Z}}$$
of a filtered complex $K$, as the graded module over $\mathbb{F}[t]$ with left multiplication by $t$ on $j$-homogeneous elements given by the linear map

$$\phi_j : H_n(K_j; \mathbb{F}) \rightarrow H_n(K_{j+1}; \mathbb{F})$$

induced by the inclusion $K_j \hookrightarrow K_{j+1}$. Since then, $PH_n(K; \mathbb{F})$ is referred to in the literature as a persistence module. More generally [6, 7], let $J$ and $C$ be categories with $J$ small (i.e., so that its objects form a set). The category of $J$-indexed persistence objects in $C$ is defined as the functor category $\text{Fun}(J, C)$; its objects are functors $F : J \rightarrow C$, and its morphisms are natural transformations $\varphi : F_1 \Rightarrow F_2$.

The typical indexing category comes from having a partially ordered set $(P, \preceq)$, and letting $P$ denote the category with objects $\text{Obj}(P) = P$, and a unique morphism from $p_1$ to $p_2$ whenever $p_1 \preceq p_2$. We’ll abuse notation and denote this morphism by $p_1 \preceq p_2$, instead of the categorical notation $p_1 \rightarrow p_2$.

It can be readily checked that if $\text{Mod}_R$ denotes the category of (left) modules over a commutative ring $R$ with unity, and $g\text{Mod}_{R[t]}$ is the category of $\mathbb{Z}$-graded modules over the polynomial ring $R[t]$, then

$$\text{Fun}(\mathbb{Z}, \text{Mod}_R) \rightarrow g\text{Mod}_{R[t]}$$

$$M, \varphi \rightarrow \bigoplus_{j \in \mathbb{Z}} M(j), \bigoplus_{j \in \mathbb{Z}} \varphi_j$$

is an equivalence of categories. On the graded $R[t]$-module side, multiplication by $t$ on $j$-homogeneous elements is given by $M(j \leq j + 1) : M(j) \rightarrow M(j + 1)$. This equivalence shows why/how the evolution of homological features in a $\mathbb{Z}$-filtered complex $K$, can be encoded as the algebraic structure of the persistence module $PH_n(K; \mathbb{F})$.

2.1. Persistence Modules and Barcodes. When $PH_n(K; \mathbb{F})$ is finitely generated as an $\mathbb{F}[t]$-module — e.g. if $K_j = \emptyset$ for $j < 0$ and $\bigcup_{j \in \mathbb{Z}} K_j$ is finite — then one has a graded isomorphism

$$PH_n(K; \mathbb{F}) \cong \bigoplus_{q=1}^{Q} t^{n_q} \cdot \mathbb{F}[t] \bigoplus \bigoplus_{l=1}^{L} (t^{m_{l\cdot}} \cdot \mathbb{F}[t])/(t^{m_{l\cdot} + d_{i,l}})$$

for $n_q, m_{l\cdot} \in \mathbb{Z}$ and $d_{i,l} \in \mathbb{N}$ [48]. The decomposition (2) is unique up to permutations, and thus the intervals

$$[n_1, \infty), [n_2, \infty), \ldots, [n_Q, \infty),$$

$$[m_1, m_1 + d_1), [m_2, m_2 + d_2), \ldots, [m_L, m_L + d_L)$$

provide a complete discrete invariant for (i.e., they uniquely determine) the $\mathbb{F}[t]$-isomorphism type of $PH_n(K; \mathbb{F})$. Moreover, this multiset recovers the barcode $\text{bcd}_n(K)$ of Edelsbrunner, Letscher and Zomorodian [27].

Carlsson and Zomorodian also observe that $PH_n(K; \mathbb{F})$ is in fact the homology of an appropriate chain complex of graded $\mathbb{F}[t]$-modules. Hence, a graded version of the Smith Normal Form [25] computes the barcode decomposition (2), providing a general-purpose algorithm. This opened the flood gates: barcodes could now be computed as a linear algebra problem for any finite filtered simplicial complex $K_0 \subset \cdots \subset K_J = K$, over any (in practice finite) field of coefficients, and up to any homological dimension. The resulting matrix reduction algorithm, implemented initially in the JPlex library (now javaPlex) [3], runs in polynomial time: its worst
time complexity is $O(N^3)$, where $N$ is the number of simplices of $K$. In fact Dmitriy Morozov exhibits in [40] a finite filtered complex of dimension 2, attaining the worst-case. This shows that the cubic bound is tight for general barcode computations.

While this sounds potentially slow, specially compared to the time complexity $O(N \cdot \alpha(N))$ of the sequential algorithm, Morozov’s example should be contrasted with filtrations arising from applications. In practice the matrices to be reduced are sparse, and computing their associated barcode decomposition takes at worst matrix-multiplication time $O(N^\omega)$ [38], where $\omega \approx 2.373$ [49]. Over the last ten years or so there has been a flurry of activity towards better implementations and faster persistent homology computations. A recent survey [42] compares several leading open source libraries for computing persistent homology. All of them implement different optimizations, exploit new theoretical developments and novel heuristics/approximations. For instance, one improvement is to first simplify the input filtered complex without changing its persistent homology (e.g., using discrete heuristics/approximations. For instance, one improvement is to first simplify the input filtered complex without changing its persistent homology (e.g., using discrete heuristics/approximations.

The shift towards algebra has had other important consequences; specifically: 1) a better understanding of stability for barcodes, and 2) several theorems describing how the choice of categories $J$ and $C$ impacts the computability of isomorphism invariants for objects in $\text{Fun}(J, C)$. Let me say a few words about stability.

### 2.2. The Stability of Persistence.

Let $X$ be a triangulable topological space (e.g., a smooth manifold) and let $f : X \to \mathbb{R}$ be a tame function (this is a generalization of being Morse). The prototypical example in TDA arises from a compact set $X \subset \mathbb{R}^d$, and letting $f_X : \mathbb{R}^d \to \mathbb{R}$ be

$$f_X(y) = \inf_{x \in X} \|x - y\|.$$ 

Thus $f_X^{-1}(-\infty, \epsilon] = X^{(\epsilon)}$. Given $f : X \to \mathbb{R}$, let $\text{bcd}_n(f)$ denote the barcode for the $n$-th persistent homology of $\{f^{-1}(\cdot, \epsilon]\}_{\epsilon \in \mathbb{R}}$. Drawing inspiration from Morse theory, Cohen-Steiner, Edelsbrunner and Harer introduced in (2007) [18] two foundational ideas: (1) the bottleneck distance $d_B(\cdot, \cdot)$ between barcodes, and (2) the stability theorem asserting that for tame $f, g : X \to \mathbb{R}$ one has that

$$d_B(\text{bcd}_n(f), \text{bcd}_n(g)) \leq \|f - g\|_\infty$$

In particular, if $X, Y \subset \mathbb{R}^d$ are compact and $X = \{X^{(\epsilon)}\}_\epsilon, Y = \{Y^{(\epsilon)}\}_\epsilon$, then

$$d_B(\text{bcd}_n(X), \text{bcd}_n(Y)) \leq d_H(X, Y).$$

This inequality implies that slight changes to the input data change the barcodes slightly, which is key for applications where (Hausdorff) noise plays a role.

Towards the end of 2008 Chazal et. al solidified the idea of stability with the introduction of interleavings for $\mathbb{R}$-indexed persistence modules [13]. The construction is as follows. For $\delta \geq 0$ let $T_\delta : \mathbb{R} \to \mathbb{R}$ be the translation functor $T_\delta(\epsilon) = \epsilon + \delta$. An $\delta$-interleaving between two persistence vector spaces $V, W : \mathbb{R} \to \text{Mod}_\mathbb{F}$ is a pair $(\varphi, \psi)$ of natural transformations

$$\varphi : V \Rightarrow W \circ T_\delta \quad \text{and} \quad \psi : W \Rightarrow V \circ T_\delta$$

so that $\psi_{\epsilon+\delta} \circ \varphi_\epsilon = V(\epsilon \leq \epsilon + 2\delta)$ and $\varphi_{\epsilon+\delta} \circ \psi_\epsilon = W(\epsilon \leq \epsilon + 2\delta)$ for all $\epsilon \in \mathbb{R}$. The interleaving distance between $V$ and $W$, denoted $d_I(V, W)$, is defined as the infimum over all $\delta \geq 0$ so that $V$ and $W$ are $\delta$-interleaved, if interleavings exist. If

\footnote{A similar result was established in [21] for $n = 0$.}
there are no interleavings, the distance is defined as $\infty$. It readily follows that $d_I$ is an extended (since infinity can be a value) pseudometric on $Fun(\mathbb{R}, \text{Mod}_\varphi)$, and that $d_I(V, W) = 0$ whenever $V \cong W$. The converse, however, is false in general (more on this later).

Chazal et. al. [13] show that if $V : \mathbb{R} \rightarrow \text{Mod}_\varphi$ satisfies $\text{rank}(V(\epsilon < \epsilon')) < \infty$ for all pairs $\epsilon < \epsilon'$, this is called being $q$-tame, then $V$ has a well-defined barcode $\text{bcd}(V)$ (see [20] for a shorter proof when $\text{dim}_\mathbb{R}V(\epsilon) < \infty$ for all $\epsilon$; this is called being pointwise finite). Moreover, if $V, W$ are $q$-tame, then one has the algebraic stability theorem $d_B(\text{bcd}(V), \text{bcd}(W)) \leq d_I(V, W)$. This turns out to be an equality:

$$d_B(\text{bcd}(V), \text{bcd}(W)) = d_I(V, W)$$

which nowadays is referred to as the Isometry Theorem; the first known proof is due to Lesnick [36].

As I said earlier, $d_I(V, W)$ can be zero for $V$ and $W$ nonisomorphic, and thus $\text{bcd}(V)$ is not a complete invariant in the $q$-tame $\mathbb{R}$-indexed case. This can be remedied as follows. Let $qFun(\mathbb{R}, \text{Mod}_\varphi)$ denote the full subcategory of $Fun(\mathbb{R}, \text{Mod}_\varphi)$ comprised of $q$-tame persistence modules. The ephemeral category $eFun(\mathbb{R}, \text{Mod}_\varphi)$, is the full subcategory of $qFun(\mathbb{R}, \text{Mod}_\varphi)$ with objects $V : \mathbb{R} \rightarrow \text{Mod}_\varphi$ satisfying $V(\epsilon < \epsilon') = 0$ for all $\epsilon < \epsilon'$. The observable category $oFun(\mathbb{R}, \text{Mod}_\varphi)$ is the quotient category

$$qFun(\mathbb{R}, \text{Mod}_\varphi)/eFun(\mathbb{R}, \text{Mod}_\varphi)$$

As shown by Chazal et. al. in [14], $d_I$ descends to an extended metric on the observable category, and taking barcodes

$$\text{bcd} : (oFun(\mathbb{R}, \text{Mod}_\varphi), d_I) \rightarrow (\text{Bcd}, d_B)$$

is an isometry. Hence, the barcode is a complete invariant for the isomorphism type of observable $\mathbb{R}$-indexed persistence vector spaces. In summary, the modern view of stability is algebraic; persistence modules are compared via interleavings, which one then tries to relate to the bottleneck distance between the associated barcodes if they exist.

### 2.3. Changing Indexing Categories: Multi-d Persistence, Quivers and Zigzags

One of the early realizations in TDA was the usefulness of having filtrations indexed by more than one parameter (1999) [30]. For instance, given a data set $X \subset \mathbb{R}^d$ one might want to focus on densely-populated regions [9], or portions with high/low curvature [12]. This leads naturally to $\mathbb{Z}^n$-filtered complexes: $\{K_u\}_{u \in \mathbb{Z}^n}$, $u = (u_1, \ldots, u_n)$, where $K_u \subset K_v$ whenever $u \preceq v$ (i.e., $u_1 \leq v_1, \ldots, u_n \leq v_n$). In this multi-filtered complex, each filtering direction $u_1, \ldots, u_n$ is meant to capture an attribute: e.g. distance/scale, density, curvature, etc. Taking homology with coefficients in $\mathbb{F}$ yields objects in $Fun(\mathbb{Z}^n, \text{Mod}_\varphi)$, and just like before, $\mathbb{Z}^n$-indexed persistence $\mathbb{F}$-vector spaces correspond to $n$-graded modules over the $n$-graded polynomial ring $P_n := \mathbb{F}[t_1, \ldots, t_n]$. Parameterizing the isomorphism classes of said modules, for $n \geq 2$, turns out to be much more involved than the barcodes from $n = 1$. Indeed, around 2009 Carlsson and Zomorodian [11] showed that the isomorphism type of a finitely generated $n$-graded $P_n$-module is uniquely determined by the following data: two finite multisets $\xi_0, \xi_1 \subset \mathbb{Z}^n$ encoding the location and multiplicity of birth-death events, and a point in the quotient of an algebraic variety $\mathcal{R}\mathcal{F}(\xi_0, \xi_1)$ by the algebraic action of an algebraic group $G_{\xi_0}$. The multisets $\xi_0, \xi_1$ are the
discrete portions of the resulting isomorphism invariant, while $\mathcal{RF}(\xi_0, \xi_1) / G_{\xi_0}$ parameterizes the (potentially) continuous part. Here is an example due to Carlsson and Zomorodian [11] illustrating how complicated this quotient can be. For $n = 2$, consider the isomorphism classes of $P_2$-modules having $\xi_0 = \{(0,0), (0,0)\}$ and $\xi_1 = \{(3,0), (2,1), (1,2), (0,3)\}$. If $\text{Gr}_1(\mathbb{P}^2)$ denotes the Grassmannian of lines in $\mathbb{P}^2$, then

$$\mathcal{RF}(\xi_0, \xi_1) = \text{Gr}_1(\mathbb{P}^2) \times \text{Gr}_1(\mathbb{P}^2) \times \text{Gr}_1(\mathbb{P}^2) \times \text{Gr}_1(\mathbb{P}^2)$$

and $G_{\xi_0}$ turns out to be the degree 2 general linear group $\text{GL}_2(\mathbb{F})$ acting diagonally on $\text{Gr}_1(\mathbb{P}^2)^4$. Since $\text{Gr}_1(\mathbb{P}^2)^4 / \text{GL}_2(\mathbb{F})$ contains a copy of $\mathbb{F} \times \{0,1\}$, and each point in this set yields a distinct isomorphism class of $P_2$-modules, it follows that there is no complete discrete invariant for (finite!) multi-d persistence.

The vast majority of recent results from multidimensional persistence focus on computable descriptors/visualizations of its intricate algebraic structure. Besides introducing the parametrization

$$\xi_0, \xi_1, \mathcal{RF}(\xi_0, \xi_1) / G_{\xi_0},$$

Carlsson and Zomorodian also propose the rank invariant: For a q-tame module $V : \mathbb{Z}^n \to \text{Mod}_\mathbb{F}$, it is defined as the function $\rho_V$ sending each pair $u \preceq v$ in $\mathbb{Z}^n$ to the integer rank $V(u \preceq v)$. $\rho_V$ is computable (see [10] for a polynomial-time algorithm), it is discrete, and an invariant of the isomorphism type of $V$. When $n = 1$ one can recover $\text{bcd}(V)$ from $\rho_V$ and viceversa, and thus $\rho_V$ is complete in the 1-dimensional case. Knudson notes in [34] that $\xi_0(V)$ and $\xi_1(V)$ are in fact the locations/multiplicities of birth events in the torsion modules $\text{Tor}_0^{P_n}(V, \mathbb{F})$ and $\text{Tor}_1^{P_n}(V, \mathbb{F})$, respectively; here $\mathbb{F}$ is identified with the $P_n$-module $\mathbb{F}[t_1, \ldots, t_n]/(t_1, \ldots, t_n)$.

The higher-dimensional analogs $\text{Tor}_j^{P_n}(V, \mathbb{F})$, $j \geq 2$, lead to a family of finite multisets $\xi_j(V) \subset \mathbb{Z}^n$, each with its own geometric interpretation, serving as isomorphism invariants for $V$. Other approaches to invariants for multidimensional persistence include the Hilbert Series of Harrington et. al. [33], the extended algebraic functions of Skrzyzaln and Carlsson [46], and the feature counting invariant of Scolamiero et. al. [45]. Lesnick and Wright have recently released RIVET, the Rank Invariant Visualization and Exploration Tool [37]. Put simply, RIVET uses the fact that if $V : \mathbb{R}^2 \to \text{Mod}_\mathbb{F}$ is q-tame and $L \subset \mathbb{R}^2$ is a line with nonnegative slope (hence a totally ordered subset of $(\mathbb{R}^2, \preceq)$), then $V^L : L \to \text{Mod}_\mathbb{F}$, the 1-dimensional persistence vector space obtained by restricting $V$ to $L$, has a well-defined barcode $\text{bc}(V^L)$. The key feature in RIVET is a graphical interface which, for finite bi-filtrations, displays $\text{bc}(V^L)$ interactively as the user varies $L$. This is particularly useful for parameter selection and the exploratory analysis of data sets with filtering functions.

Multidimensional persistence is a great example of how a seemingly innocuous change in indexing category, say from $\mathbb{Z}$ to $\mathbb{Z}^2$, can lead to a widely different and much more complicated classification problem. With this in mind, one would like to have a systematic approach to address the ensuing complexity. The representation theory of Quivers [24] offers one such avenue. It turns out that the classification of finite J-indexed persistence vector spaces $V : J \to \text{Mod}_\mathbb{F}$ can be studied directly from the shape of the indexing category $J$. Indeed, let $G(J)$ be the finite directed (multi)graph with the objects of $J$ as vertices, and one arrow for every morphism...
that is neither an identity nor a composition. Also, let $\tilde{G}(J)$ be the undirected graph obtained from $G(J)$ by forgetting arrow directions. When $\tilde{G}(J)$ is acyclic, Gabriel’s theorem [31] implies that pointwise finite objects in $\text{Fun}(J, \text{Mod}_F)$ can be classified via complete discrete invariants, if and only if the connected components of $\tilde{G}(J)$ are Dynkin diagrams of the types described in Figure 4 below.

Here is an example of how this result can be used to avoid unpleasant surprises: Suppose that $G(J)$ is the graph with vertices $x_0, \ldots, x_N$ and $N \geq 5$ edges $x_n \to x_0$, $n = 1, \ldots, N$ (see Examples 3 and 8 in [24]). While the resulting $J$-indexed persistence vector spaces $V : J \to \text{Mod}_F$ may look simple (just star-shaped, right?), the connected graph $\tilde{G}(J)$ is not a Dynkin diagram, and the ensuing classification problem is in fact of “wild type”: complete invariants must include continuous high-dimensional pieces, just like in multidimensional persistence.

These ideas entered the TDA lexicon around 2010 with the definition of Zigzag persistence by Carlsson and de Silva [8]. Regular persistence addresses the problem of identifying stable homological features in a monotone system of spaces and continuous maps

$$X_1 \to X_2 \to \cdots \to X_J.$$

Zigzag persistence, on the other hand, is a generalization to the non-monotone case. Here is a practical example: suppose one has an ordered sequence of spaces $X_1, \ldots, X_J$ (e.g., from time varying data), but no obvious maps $X_j \to X_{j+1}$. The need to track topological features as $j$ varies leads one to consider the system

$$X_1 \leftarrow X_1 \cup X_2 \leftarrow X_2 \leftarrow \cdots \leftarrow X_j \cup X_{j+1} \leftarrow \cdots \leftarrow X_J$$

and the resulting zigzag diagram

$$V_1 \to V_2 \to \cdots \to V_n$$

at the homology level. More generally, a (finite) zigzag is a sequence of vector spaces $V_1, \ldots, V_n$ and linear maps $V_j \to V_{j+1}$ or $V_j \leftarrow V_{j+1}$. The sequence of arrow directions, e.g. $\tau = (\text{left}, \text{left}, \text{right}, \ldots, \text{right}, \text{left})$, is the zigzag type. Since in this case any choice of $\tau$ forces the indexing category $J_\tau$ to satisfy $\tilde{G}(J_\tau) = A_n$ (one of the aforementioned Dynkin diagrams), then Gabriel’s theorem implies that finite zigzags

$$V : J_\tau \to \text{Mod}_F$$
are completely classified by a discrete invariant. Just as for regular 1-dimensional persistence the invariant turns out to be a barcode, which can be efficiently computed \cite{38}, and for which there is a zigzag stability theorem \cite{5} recently established by Botnan and Lesnick.

When the graph $\tilde{G}(J)$ has cycles, the functoriality of objects in $\text{Fun}(J, \text{Mod}_F)$ is captured by the notion of a quiver with relations. The taxonomy from Gabriel’s theorem no longer applies, but one can still find some answers in the representation theory of associative algebras. A particularly important instance is when the cycles of $\tilde{G}(J)$ are not oriented cycles in $G(J)$; in this case the algebras of interest are finite dimensional (hence Artinian) and Auslander-Riten theory \cite{4} becomes relevant.

Escolar and Hiraoka \cite{28} have recently put these ideas to use in the context of persistent objects indexed by commutative ladders; that is, the persistence of a morphism between two zigzags of the same type:

\[
\begin{array}{cccccccc}
\bullet & \longrightarrow & \bullet & \longleftrightarrow & \cdots & \longrightarrow & \bullet & \longrightarrow & \bullet & \longrightarrow \\
\downarrow & & \downarrow & & & & \downarrow & & \downarrow & \downarrow & \downarrow \\
\bullet & \longrightarrow & \bullet & \longleftrightarrow & \cdots & \longrightarrow & \bullet & \longrightarrow & \bullet & \longrightarrow
\end{array}
\]

The resulting theory sits somewhere between zigzag persistence and multi-dimensional persistence: short ladders (length $\leq 4$) have complete discrete invariants, but longer ones do not. Escolar and Hiraoka present an algorithm for computing these invariants, and also an interesting application to computational chemistry.

I think this is a good place for me to stop; hopefully it is also a good starting point for the reader interested in persistent homology. There are several books covering many of the ideas I presented here, as well as many others. The interested reader would certainly benefit from these resources \cite{32, 15, 26, 43}.

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