Granular computing and reasoning based on rough logic

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Abstract: This paper makes an effective combination of rough set and granular computing,
and proposes a logic inference form based on rough granular computing. The advantage of
rough logic granular computing lies in its demonstration form which is given by array, it can
illustrate the formula semantic more directly, and furthermore the reasoning results can be got
according to the data information searched momentarily.

Chinese classification code: O159

1. Introduction

The expressive method used by granular computing[1][2] which is based on rough set[3][4][5][6][7] is very
special. Each granule is expressed in a binary set. The first element of the set is a logic formula, and
the second element is a semantic set corresponding to the formula.

Because the two elements of this binary set make up a cohesive whole that shows both syntax and
semantic set, so it is called granule. It is both logical and set theory.

This makes it possible to use both logical method and set theory method in deductive or other
approximate reasoning[8]. So this is also the advantage of granular computing.

2. Basic Theory of Rough Logic

Definition 1.1[9] Suppose that IS = (U, A) is a information system. U is all of the objects discussed.
A is the property set. Define (a, v) or a_v is an atomic formula in the rough logic, where a ∈ A
and v ∈ V_a is the attribute value. The formula which combines the atomic formula with the classical
logical conjunction is called a compound formula.

Definition 1.2[9] Rough logic (RL) is composed of attribute set A, attribute values set
V = ∪_{a∈A}V_a, usual logical connective word, parentheses and the following formula

(1)All the atomic formula of form a_v is the wff in RL. a ∈ A, v ∈ V_a;

(2)Suppose that ϕ and ψ are wffs in RL, then ¬ϕ, ϕ ∨ ψ, ϕ ∧ ψ, ϕ → ψ, ϕ ↔ ψ are wffs
in RL;

(3)The formula which is got from finite number of repeated using (1) and (2) is wffs in RL.

Definition 1.3[9] Giving information system IS = (U, A), suppose that F(●) is a Rough logic
formula in IS, define formula F(●):m(F(●)) = {u ∈ U | u |≈_{IS} F(●)\}.
Definition 1.4 Suppose that $F_1, F_2 \in RL_{IS}$ is the Rough logic formula on the given information system, the semantic sets are $m(F_1), m(F_2)$, then it gets:

1. $m(\neg F_1) = U - T_{lu}(F_1)$;
2. $m(F_1 \lor F_2) = m(F_1) \cup m(F_2)$;
3. $m(F_1 \land F_2) = m(F_1) \cap m(F_2)$;
4. $m(F_1 \rightarrow F_2) = m(\neg F_1) \cup m(F_2)$;
5. $m(F_1 \leftrightarrow F_2) = (m(\neg F_1) \cup m(F_2)) \cap (m(\neg F_2) \cup m(F_1))$.

With rough logical formula, the truth value of rough logical formula is defined as follows:

Definition 1.5 Suppose that $F_1, F_2 \in RL_{IS}$ is the rough logic formula in IS, then the true value is defined as follows:

(1) $T_{lu}(F_1) = K(m(F_1)) / K(U)$;
(2) $T_{lu}(\neg F_1) = 1 - T_{lu}(F_1)$;
(3) $T_{lu}(F_1 \lor F_2) = T_{lu}(F_1) \lor T_{lu}(F_2)$;
(4) $T_{lu}(F_1 \land F_2) = T_{lu}(F_1) \land T_{lu}(F_2)$;
(5) $T_{lu}(\forall x)F_1(x) = \min\{T_{lu}(F_1(e_1)), T_{lu}(F_1(e_2)), \ldots, T_{lu}(F_1(e_n))\}$.

3. The theory of Granular computing on rough logic

Definition 2.1 Suppose that $IS = (U, A)$ is a information system, $\phi$ is the rough logic formula in $IS$, $m(\phi)$ is its semantic set, define binary pair $G = (\phi, m(\phi))$ as the basic Granular about $\phi$ in $IS$.

All the Granular in $IS$ is $L_{IS}$.

If $m(\phi) = \emptyset$, then can call $\phi$ is false in $IS$, it can be written as $\nexists IS G$;
If $m(\phi) = U$, then can call $\phi$ is true in $IS$, it can be written as $\exists IS G$;
If $\emptyset \subseteq m(\phi) \subseteq U$, then can call $\phi$ is satisfied in $IS$, it can be written as $\sar IS G$;
If $m_s(\phi) = \{x \in U : x \sar IS, \phi = B_s(m(\phi))\}$, Where $m_s$ is the lower semantic function of $\phi$, then $\phi$ can be called rough lower true in $IS$, it can be written as $\lra IS G$;
If $m_s(\phi) = \{x \in U : x \sar IS, \phi = \emptyset\}$, then $\phi$ can be called rough lower false in $IS$, it can be written as $\lnot IS \phi$;
If $m_s(\phi) = \{x \in U : x \sar IS, \phi = B_s(m(\phi))\}$, Where $m_s$ is the lower semantic function of $\phi$, then $\phi$ can be called rough upper true in $IS$, it can be written as $\lra IS \phi$;
If $m_s(\phi) = \{x \in U : x \sar IS, \phi = \emptyset\}$, then $\phi$ can be called rough upper false in $IS$, it can be written as $\lnot IS \phi$;
If $\emptyset \subseteq m_s(\phi) \subseteq B_s(m(\phi)) \subseteq U$, then $\phi$ can be called rough satisfied in $IS$, it can be written as $\sar IS \phi$;
If $m_s(\phi) = \{x \in U : x \sar IS, \phi = U\}$, then $\phi$ can be called rough true in $IS$, it can be written as $\lra IS \phi$;
If $m_s(\phi) = \{x \in U : x \sar IS, \phi = \emptyset\}$ and $m_s(\phi) = \{x \in U : x \sar IS, \phi = \emptyset\} = \emptyset$, then $\phi$ can be
called rough incompatible in \( IS \).

For decision rules \( \phi \to \varphi \) in a decision system, it can be written as \((\phi, \varphi)\). Its semantic set is written as \((m(\phi), m(\varphi))\). The basic Grain of the rule \( \phi \to \varphi \) is defined as \(((\phi, m(\phi)), (\varphi, m(\varphi)))\).

**Definition 2.2** Suppose that \( G = (\phi, m(\phi)), G' = (\varphi, m(\varphi)) \) are two Grain of statements. Their operation definitions of connectives \( \neg \) (negation), \( \oplus \) (or), \( \otimes \) (and), \( \Rightarrow \) (implication) \( \Leftrightarrow \) (equivalence) are as follows:

1. \( - (\phi, m(\phi)) = (- \phi, U - m(\phi)) \);
2. \( (\phi, m(\phi)) \oplus (\varphi, m(\varphi)) = (\phi \lor m(\phi) \cup m(\varphi)) \);
3. \( (\phi, m(\phi)) \otimes (\varphi, m(\varphi)) = (\phi \land m(\phi) \cap m(\varphi)) \);
4. \( (\varphi, m(\varphi)) \Rightarrow (\phi, m(\phi)) = (\phi \to m(\varphi) \subseteq m(\phi)) \lor ((m(\varphi) \subseteq m(\phi) \land m(\varphi) \subseteq m(\phi))) \)
5. \( (\varphi, m(\varphi)) \Leftrightarrow (\phi, m(\phi)) = (\phi \leftrightarrow m(\varphi) \subseteq m(\phi) \land m(\varphi) \subseteq m(\phi)) \)

**4. Semantic Reasoning Based On Rough Granular**

Based on the grammar and semantic of granular computing, the semantic reasoning of granular is given.

**Theorem 3.1** Suppose that \( G = (\phi, m(\phi)), G' = (\varphi, m(\varphi)) \in L_{IS} \), then: \( \models_r (G \Rightarrow G') \otimes G \), then \( \models_r G' \).

Proof: Because \( \models_r (G \Rightarrow G') \otimes G \), so it gets
\[
B^*(m(\phi \to \varphi)) \subseteq B^*(m(\phi \to \varphi)) \cap m(\phi) \subseteq B^*(m(\phi \to \varphi)) \cap B^*(m(\phi))
\]
\[
B^*(m(\phi \to \varphi)) \subseteq B^*(m(\phi \to \varphi)) \cap \neg \phi \cup \varphi \subseteq B^*(m(\phi \to \varphi)) \cap B^*(\neg \phi) \cup B^*(\varphi)
\]
\[
B^*(m(\phi)) = U
\]
then \( \models_r G' \).

**Theorem 3.2** Suppose that \( G = (\phi, m(\phi)), G' = (\varphi, m(\varphi)) \in L_{IS} \), if \( \models_r G \Rightarrow G' \) and \( \models_{IS} G \), then \( \models_r G' \).

Proof: Because \( \models_r G \Rightarrow G' \), so it gets
\[
B^*(m(\phi \to \varphi)) = B^*(m(\phi \to \varphi)) \cup B^*(m(\phi)) = U
\]
and because \( \models_{IS} G \), so \( m(\phi) = U \), \( B^*(m(\phi)) = U - B_*(m(\phi)) = \emptyset \), \( B^*(m(\phi)) = U \).

**Theorem 3.3** Suppose that \( G = (\phi, m(\phi)), G' = (\varphi, m(\varphi)) \) are two granular statements on rough granular. Then there are two conclusions:

(1) If \( \models_r G \Rightarrow G' \), and \( m(\phi) \cup m(\varphi) = U \), then \( \models_r G' \);
(2) If \( \models_r G \), and \( m(\phi) \cap m(\varphi) = \emptyset \), then \( \not\models_r G \Rightarrow G' \).

Proof: (1) Because \( \not\models_r G \Rightarrow G' \), so
\[
B^*(m(\phi \to \varphi)) = B^*(m(\phi \to \varphi)) \cup B^*(m(\phi)) = U
\]
And \( m(\phi) \cup m(\varphi) = U \), then \( U - m(\phi) \subseteq m(\varphi) \)
\[
B^*(m(\phi \to \varphi)) \subseteq B^*(m(\phi)) = U
\]
It gets the conclusion.
(2) Because $|_{R}G$, so $B^*(m(\phi)) = U$
And because $m(\phi) \cap m(\phi) = U$, so $m(\phi) \subseteq U - m(\phi)$, then
$$B_*(m(\phi)) \subseteq B_*(U - m(\phi)) = U - B^*(m(\phi))$$
So gets the conclusion
$$B_*(m(\phi \to \phi)) = B_*(m(\neg\phi \lor \phi)) = B_*(m(\neg\phi) \cup m(\phi)) = B_*(U - m(\phi)) = U - B^*(m(\phi)) = \emptyset$$

**Theorem 3.4** There are Granular $G = (\phi, m(\phi)) \in L_{IS}$ in approximate space, which makes $G$ and $\neg G$ are true in $IS$.
Proof: If $G$ is rough and not compatible in $IS$, then $B^*(m(\phi)) = U$, and $B_*(m(\phi)) = \emptyset$. So, $B^*(m(\neg\phi)) = B^*(U - m(\phi)) = U - B_*(m(\phi)) = U$.

**Theorem 3.5** For Granular $G = (\phi, m(\phi)) \in L_{IS}$ in approximate space $IS = (U, A)$, If $G$ is rough and not compatible in $IS$, then $\neg G$ is rough and not compatible in $IS$.
Proof: Because $G$ is rough and not compatible in $IS$, then $B^*(m(\phi)) = U$, and $B_*(m(\phi)) = \emptyset$ so,
$$B^*(m(\neg\phi)) = B^*(U - m(\phi)) = U - B_*(m(\phi)) = U$$
$$B_*(m(\neg\phi)) = B_*(U - m(\phi)) = U - B^*(m(\phi)) = \emptyset$$
Then can get the conclusion that $G$ is rough and not compatible in $IS$.

**Theorem 3.6** For Granular $G = (\phi, m(\phi)) \in L_{IS}$ in approximate space $IS = (U, A)$.
(1) If $G$ is rough and not compatible in $IS$, then for all $G' = (\phi, m(\phi)) \in L_{IS}$, there are $|_{R}G \oplus G'$ and $|_{R}G \otimes G'$, namely $G \oplus G'$ is rough and true in $IS$, $G \otimes G'$ is rough and false in $IS$.
(2) There is $G' = (\phi, m(\phi)) \in L_{IS}$, which makes $|_{R}G \oplus G'$ and $|_{R}G \otimes G'$, namely $G \oplus G'$ is rough and true in $IS$, $G \otimes G'$ is rough and false in $IS$.
Proof: (1) Because $G$ is rough and not compatible in $IS$, then $B^*(m(\phi)) = U$, and $B_*(m(\phi)) = \emptyset$. For all the rough Granular $G' = (\phi, m(\phi))$,
$$B^*(m(\phi \lor \phi)) = B^*(m(\phi) \cup m(\phi)) = B^*(m(\phi)) \cup B^*(m(\phi)) = U$$
$$B_*(m(\phi \land \phi)) = B_*(m(\phi) \cap m(\phi)) \subseteq B_*(m(\phi)) \cap B_*(m(\phi)) = \emptyset$$
then, $B_*(m(\phi \land \phi)) = \emptyset$
(2) For rough Granular $G = (\phi, m(\phi)) \in L_{IS}$, if $G' = (\phi, m(\phi)) \in L_{IS}$, and satisfies
$$m(\phi) \cup m(\phi) = U \text{ and } m(\phi) \cap m(\phi) = \emptyset$$
Then
$$B^*(m(\phi \lor \phi)) = B^*(m(\phi) \cup m(\phi)) = B^*(m(\phi)) \cup B^*(m(\phi)) = U$$
$$B_*(m(\phi \land \phi)) = B_*(m(\phi) \cap m(\phi)) \subseteq B_*(m(\phi)) \cap B_*(m(\phi)) = \emptyset$$

5. Conclusion
This paper mainly discusses rough Granular logic, defines several kinds of Granular operations similar to the classical propositional logic operations, and realizes the deductive reasoning using the method of Granular calculation. The advantages of this reasoning are both logical and collective, so that the formula in each step and its evolution process of semantic set can be intuitive to see in the process of reasoning, then enhance the transparency and generalized of the reasoning.
Acknowledgement
Project supported by the National Natural Science Foundation of China (No.11732010, 11572221, 11872272) and the Scientific Research Program Project of Tianjin Education Commission (No.2018KJ267).

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