Load-Aware Energy Efficient Adaptive Large Scale Antenna System

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ABSTRACT

This paper proposes an adaptive large scale antenna system (ALSAS) for enhancing energy efficiency in low density wireless network scenarios. The proposed ALSAS comprises of two stages, a novel adaptive discontinuous transmission (ADTx) stage and an antenna array optimization (AAO) one. The basic idea is to utilize prior knowledge of the users’ quality of service (QoS) requirements as well as precoding selection in the ADTx stage to maximize the transmitter hibernation periods subject to a certain complexity constraint. In the AAO stage, further power saving is achieved by reducing the number of active antenna elements subject to a certain QoS requirement. It is shown that, relative to conventional large scale antenna system (LSAS), the proposed ALSAS system achieves significant energy efficiency improvements under various scenarios. The results show that the proposed technique can provide energy efficiency improvement between 125% and 1124% in the suburban scenario, and between 196% and 952% in the rural scenario. It is also demonstrated that for rural environments with relatively small short inter-site-distance (ISD) values, ALSAS can provide up to 500% power saving for the fixed bit rate requirement case.

INDEX TERMS

Massive MIMO, discontinuous transmission, adaptive precoding energy efficiency, QoS, suburban and rural scenarios.

I. INTRODUCTION

It is well-known that multiple-input multiple-output (MIMO) technology can significantly increase the network capacity through spatial multiplexing [1]. The fifth generation (5G) New Radio (NR) standard supports a single user equipment (UE) in single-user MIMO with a maximum of eight antennas for the down link (DL) and four for the uplink (UL) [2, pp. 29]. It also supports multi-antenna data transmission for multiple UEs, multi-user MIMO, with a maximum of twelve transmission layers for the DL and UL. MIMO, in the form of large scale antenna system (LSAS), or massive MIMO, is used as a multiuser multiplexing technique for its ability to satisfy the growing demands in wireless networks. In fact, having a large number of antennas does not only allow increasing the number of multiplexed users, but it also reduces the channel variation across the spectrum making the channel response flat which could increase the achievable capacity [3], [4]. Furthermore, beamforming in LSAS reduces the interference for the cell-edge users and allows full spectrum utilization without soft frequency reuse as in the Long Term Evolution (LTE) standard [5]. Another advantage of LSAS is the reduced isotropic radiated power (EIRP) per antenna, which simplifies the antenna chain design and lowers the power consumption per antenna chain [6]. That is, each LSAS antenna can employ a radio frequency (RF) chain similar to that of a small cell with a more efficient power amplifier (PA), orthogonal frequency division multiplexing (OFDM) modulator, low noise amplifier (LNA), etc. For instance, a small cell PA can have 50% power efficiency compared to 25% in a conventional macro base station (MBS) [7]. Several practical demonstrations have confirmed the tremendous advantage of LSAS. For example, Samsung’s 64 × 64 massive MIMO test revealed a 640% peak throughput improvement compared to a 4 × 4 MIMO [8].
The combination of low power consumption and high achievable throughput is what makes LSAS an ideal technology for 5G networks. This is especially true for a high user-density scenario where the number of multiplexed users is large. On the other hand, in low density environments, such as off-peak urban, suburban and rural scenarios, a large number of antennas may become counterproductive [7], [9]. Therefore, energy efficiency (EE) improvement will lead to a significant power reduction [7]. One way to achieve high EE using LSAS is by taking into account the general behavior of the UEs. As shown in [10], more than 40% of mobile cellular traffic is used for low data rates such as e-mail, short messages and voice over Internet protocol (VoIP), which usually require a bit rate of less than 1.5 Mbps. Although the remaining 60% of the traffic corresponds to video, most live video streaming applications require less than 0.5 Mbps for both watching and broadcasting [11]. Therefore, a full buffer scenario is not common for most cases, making discontinuous transmission (DTx) [12] an effective tool to improve EE in low density LSAS networks. DTx was originally used to reduce interference in communication channels [13], and later for reducing power consumption in LTE [14]–[19], by regularly putting a certain base station in hibernation mode. It is shown in [14] that significant power saving can be achieved when the traffic load is low. Most existing work on DTx considers wireless systems with a small number of antennas, and employ interference and transmission coordination [20]. To the best of the authors’ knowledge, DTx has never been incorporated into LSAS. The closest implementation of DTx in a system with a large number of antennas is in massive distributed antenna system [21].

In addition to DTx, extensive work has been invested to optimize the EE in LSAS. In general, such metric can be optimized using well-known techniques such as the Dinkelbach and fractional programming algorithms [22], [23]. For example, Fung et al. [24] proposed a solution to a non-convex optimization problem with nonorthogonal multiple access (NOMA) by using constraints that include the outage probability. In [25] the authors addressed maximizing the EE subject to specific power and data rate constraints, while in [26] the authors claimed to achieve near-optimal EE using joint antenna selection, power optimization and user selection. In [27] a low-complexity EE optimization solution based on maximum transmitted power and minimum data rate constraint is proposed based on fractional programming, learning, and game theory. Other works focused on the user association and/or power allocation such as [28], [29].

Given the aforementioned discussion, reducing the power consumption in wireless networks under low density scenarios is indispensable. Therefore, by capitalizing on our work [30]–[32], we propose an adaptive LSAS (ALSAS) technique that utilizes a combination of a novel adaptive DTx (ADTx) combined with precoding selection. The proposed ADTx exploits the quality of service (QoS) information, such as bit rate and latency requirements, and the flexibility of switching between different precoding techniques to maximize the antennas hibernation periods. Although there are many precoding techniques to choose from, such as conjugate beamforming (CJ), zero forcing (ZF), minimum mean square error (MMSE), and dirty paper precoding [33], [34], we confine the choice here to ZF and CJ, which are prominent precoding techniques [4]. It is worth noting that ZF has lower complexity than MMSE while can also provide near-optimal sum rate when the number of antennas is much larger than the number of users [35], pp. 67], which is the case in this paper. While CJ offers low complexity at the expense of some performance degradation when the number of users is relatively large, ZF trades complexity for improved performance when the number of users is large. After the ADTx stage, ALSAS performs antenna array optimization (AAO) to increase the hibernation period and achieve additional transmission power reduction. It will be shown that the proposed technique significantly improves EE in LSAS systems, especially when the network load is low, in both suburban and rural scenarios, while satisfying the QoS. Furthermore, the proposed technique provides a significant EE improvement when the cell size is reduced. To validate our claims, mathematical analysis corroborated by simulation results are presented for several use cases of interest. As compared to [30]–[32], this paper extends and refines the DTX scheme to a full adaptive one and tailors its optimization to low density sub-urban and rural scenarios. Additionally, it includes both analytical and simulation results and extends the results and discussion sections to include various parameters of interests such as spectral efficiency and outage probability, in addition to the EE enhancement achieved in different environments.

The rest of the paper is organized as follows. First, the system model and basic concepts are described in Section II. The proposed ALSAS is explained in Section III. ADTx is analyzed and evaluated in Section IV. Numerical results and conclusions are drawn in Sections V and VI, respectively.

A. SYMBOLS AND NOTATIONS

The following notations are used throughout this paper:

- Bold Upper cases and Greek letters with a bar such as \( \bar{\mathbf{A}} \) and \( \bar{\boldsymbol{A}} \), lower case and Greek letters such as \( \mathbf{a} \) and \( \boldsymbol{a} \) denote vectors (sets) with elements \( A_i \), \( \bar{A}_i \), \( a_i \), and \( \bar{a}_i \), respectively.
- Bold upper case letters such as \( \mathbf{B} \) denote matrices with elements \( B_{m,n} \), where \( m \) denotes the row number and \( n \) denotes the column number. The \( l \)th row or column in \( \mathbf{B} \) is denoted as \( \mathbf{B}_l \).
- The unit step function \( U(a,b) = 1 \) if \( b > a \), and 0 otherwise.
- \( [x] \), \( \lfloor x \rfloor \) and \( \lceil x \rceil \) denote the round, ceiling and floor functions, respectively.
- Subscript \( G \) and \( Z \) indicate the CJ and ZF precoding techniques, respectively.
- \( l \) is the scheduled transmission index.
- \( k \) is the user index.
- \((\cdot)^T\) is the transpose.
**TABLE 1.** List of abbreviations in alphabetical order.

| Abbreviation | Description                  |
|--------------|------------------------------|
| ADTx         | Adaptive Discontinuous Transmission |
| AAO          | Antenna Array Optimization   |
| CJ           | Conjugate Beamforming        |
| DTx          | Discontinuous Transmission   |
| EE           | Energy Efficiency            |
| PO           | Precoder Observation         |
| RT           | Release Time                 |
| RB           | Resource Block               |
| ST           | Scheduled Transmission       |
| SE           | Spectral Efficiency          |
| ZF           | Zero Forcing                 |

**TABLE 2.** List of symbols.

| Sym. | Description                        | Sym. | Description                        |
|------|------------------------------------|------|------------------------------------|
| α    | Pathloss                           | Y    | Latency duty cycle requirement     |
| P_{Tr} | Transmitted power                   | K    | Total number of users               |
| TA   | Active transmission period per frame (sec.) | R_{max} | Achievable bit rate for each OFDM symbol |
| F    | Number of OFDM symbols per frame   | A    | Allocated RBs matrix               |
| M    | Total number of Tx antennas        | L    | Total STs                          |
| M_{1} | Number of instantaneous antennas for the ADTx | ℓ | Pilot sequence length              |
| ψ    | SINR per antenna                   | N    | Number of allocated RBs           |
| γ    | SINR using precoding               | S    | Available OFDM symbols             |
| R    | Bit rate requirement               | K    | Maximum number of multiplexed UEs  |
| ε    | Energy efficiency                  | S    | Allocated symbols for the ADTx     |
| U    | Unit-step function                 | I    | Interference                       |
| T    | Set of integer numbers             | R    | Set of real numbers                |
| B    | Set of Binary numbers              | B    | Total number of RBs per OFDM symbol |
| Q    | Number of subcarriers per RB       | P_{f} | Precoding indicator               |
| C    | Set of complex numbers             |      |                                    |

Furthermore, a list of important abbreviations and symbols used throughout this paper are summarized in Tables 1 and 2, respectively.

**II. SYSTEM MODEL AND FUNDAMENTALS**

**A. SYSTEM MODEL**

This work considers a typical OFDM-based LSAS [30], [31]. To simplify the presentation, we initially consider a single MBS and later extend the model to a multi cell scenario. The considered MBS is equipped with $M$ transmit antennas, and simultaneously serves $K$ single-antenna UEs, where $M \gg K$.

The symbol vector $\mathbf{d}^n = [d_1, d_2, \ldots, d_K]^T$, $d_k$ is the data symbol of the $k$th user selected uniformly from a particular constellation such as quadrature amplitude modulation (QAM), and the index $n$ specifies the corresponding $n$th subcarrier that will be used to transmit the data vector $\mathbf{d}^n$. For linearly precoded LSAS, the vector $\mathbf{d}^n$ is multiplied by an $M \times K$ precoding matrix $\mathbf{A}^n$ to generate the precoded vector $\mathbf{v}^n$. That is,

$$\mathbf{v}^n = \mathbf{A}^n \mathbf{d}^n$$  \hspace{1cm} (1)

where $\mathbf{v}^n \in \mathbb{C}^{M \times 1}$. To generate $M$ OFDM symbols with $N$ subcarriers each, the following matrix is formed,

$$\mathbf{V} = \begin{bmatrix} \mathbf{v}_1^T \\ \mathbf{v}_2^T \\ \vdots \\ \mathbf{v}_N^T \end{bmatrix}$$  \hspace{1cm} (2)

where $\mathbf{V} \in \mathbb{C}^{N \times M}$. Consequently, each column in $\mathbf{V}$ is converted to an OFDM symbol by computing the inverse fast Fourier transform (IFFT) and appending a cyclic prefix (CP). Then each OFDM symbol is mapped to one of the $M$ antennas.

At the receiving side, each UE removes the CP samples and computes the FFT of the received sequence to produce the vector $\mathbf{r}$, which contains the information symbols of the $K$ users. At the $n$th subcarrier, $\mathbf{r}^n$ can be written as

$$\mathbf{r}^n = \mathbf{H}^n \mathbf{v}^n + \mathbf{w}^n$$ \hspace{1cm} (3)

where $\mathbf{H}^n \in \mathbb{C}^{K \times M}$ is the channel matrix and $\mathbf{w}^n \in \mathbb{C}^{K \times 1}$ represents the receiver noise that has independent and identically distributed (i.i.d.) circularly symmetric complex Gaussian entries with zero-mean and $\sigma_w^2$ variance.

The $K$ UEs require a set of different bit rates, $\mathbf{R} \in \mathbb{R}^{1 \times K}$, which are subject to certain latency requirements defined by $\mathbf{V} \in \mathbb{R}^{1 \times K}$. In this paper, latency is defined as the data transmission duty cycle for a given UE. This is defined as the ratio between the required OFDM symbols and the number of transmission slots per frame.

A modified time-division duplex (TDD) system is considered with low density UE distribution, as in rural and suburban areas. The considered frame structure follows the format described in [36], where a frame consists of $F$ OFDM data symbols and includes $\ell$ OFDM pilot symbols. Therefore, each frame has $F - \ell$ data symbols for uplink (UL) and downlink (DL) transmissions, denoted as $F_D$ and $F_U$, respectively. In this work we consider that $F_D = F_U$ and $\ell = F/10$, where $F = T_u T_{fr} / (T_g T_s)$ [36]. $T_g$ is the OFDM cyclic prefix interval, $T_s$ is the OFDM symbol interval, $T_{fr}$ is the total frame period and $T_u$ is the OFDM usable interval, $T_u = T_s - T_g$. The pilot symbols are used to estimate the channel response using the MMSE criterion [37].

As proposed by the Greentouch consortium [7], we consider separate pilot and control signal transmission to allow each MBS to hibernate when there is no data or pilot signals to be transmitted. Furthermore, it is assumed that each MBS has the ability to control the number of active antennas, which can be activated/deactivated instantly. A reuse factor of 7 is assumed for the pilot signal transmission and the pilot
signals are orthogonal in the time domain. Therefore, the pilot length for the \(k^{th}\) UE is given by \(K ≤ \ell_k ≤ \ell/7\), where \(K\) is the maximum number of instantaneous multiplexed UEs. As reported in [36], this pattern reduces pilot contamination to a value that is much lower than the noise level, making it plausible to ignore the impact of pilot contamination. In the frequency domain, the subcarriers are divided into resource blocks (RBs) each of which has 12 subcarriers, and time-spectrum will be allocated for each RB and OFDM symbol.

\[ ϱ_k = \frac{α_k^{-1} P_D}{η_k} \]  

(4)

where \(α_k\) is the pathloss coefficient between the MBS and \(k^{th}\) UE, \(P_D\) is the transmit power per antenna at the MBS and \(η_k\) is the noise power of the UE. The effective SINR of the \(k^{th}\) UE with \(M\) active antennas, \(K\) active UEs, with imperfect channel estimates, can be expressed as [4],

\[ γ_k = \frac{M K ^{-1} α_k + μ}{K} \]  

(5)

where \(μ = 0\) and \(-α_k\) for CJ and ZF precoding, respectively, and

\[ a_k = \begin{cases} \frac{ψ_k Q_k}{(ψ_k + 1) (Q_k + 1)}, & \text{CJ precoding} \\ ψ_k Q_k, & \text{ZF precoding} \end{cases} \]  

(6)

Furthermore, \(Q_k = \ell_k P_U / (α_k ν_0)\), where \(ν_0\) is the noise power at the MBS. Since a large number of antennas can eliminate small-scale fading variations [3], the effective SINR for the precoded system is the same across all RBs. Given that each OFDM symbol has a maximum of \(Ω\) RBs that can be allocated to a particular user, the maximum achievable bit rate for that user is given by,

\[ R_k^{max} = \frac{1}{T_{fr}} F Ω Q \log_2 (1 + γ_k) \text{bit/s} \]  

(7)

where \(Q\) is the number of subcarriers per RB. Based on (4)-(7), LSAS will have a set of pathloss coefficients, \(α \in \mathbb{R}^{1×K}\), received interference power set, \(I \in \mathbb{R}^{1×K}\), to calculate the SINR set \(γ \in \mathbb{R}^{1×K}\) and the achievable bit rate \(R_k^{max} \in \mathbb{R}^{1×K}\). To simplify notations, \(ψ_G, a_G, μ_G\) and \(R_k^{max}\) are used for CJ precoding and \(ψ_Z, a_Z, μ_Z\) and \(R_Z^{max}\) for ZF precoding.

\section{C. ENERGY EFFICIENCY MODEL}

The EE computed over a given frame is computed in this work in terms of the energy productivity of the system, which is defined as

\[ E = \frac{B_{fr}^κ}{P_{Efr} T_{fr}} \text{ bits/Joule} \]  

\[ = \frac{B_{fr}^κ}{P_{Efr} T_{fr}} \]  

(8)

where \(P_{Efr}\) is the total consumed power and \(B_{fr}^κ\) is the total number of bits transmitted per frame. By noting that \(B_{fr}^κ/ T_{fr}\) simply represents the total DL bit rate \(R^κ\), then

\[ E = \frac{R^κ}{P_{Efr}} \]  

\[ = \frac{1}{P_{Efr}} \sum_{k=1}^{K} \min \{ R_k, R_k^{max} \}. \]  

(9)

The data rate for each user is selected as \(\min \{ R_k, R_k^{max} \}\) to minimize the energy consumption. Therefore, the MBS will limit the total bit rate to the desired rate even when higher bit rates can be achieved.

The total power consumption \(P_{Efr}\) is related to the DL transmission model for the 2020 scenarios [7] is given by

\[ P_{Efr} = P_A + P_B + P_C + P_H + P_J \]  

(10)

where \(P_A\) denotes the power consumption of the power amplifier (PA), \(P_B\) is the digital baseband power, \(P_C\) is the antenna analog circuit power, \(P_H\) is the power for signaling and control bits, and \(P_J\) is the power consumption related to the control, backhaul and network signaling, which is typically 300\(\text{mW}\) [7], [38], [39]. The powers \(P_A, P_B\) and \(P_C\) are assumed to be adaptive where it switches between active (on) and hibernation (off) states, where at a given instant, \(M\) antennas are active for \(\bar{e} \in \mathbb{R}^{1×M}\) percent of the time slot. Hence, \(P_A = \sum_{m=1}^{M} \bar{e}_m P_A^\text{on} + (1 - \bar{e}_m) P_A^\text{off} = \sum_{m=1}^{M} \bar{e}_m P_A^\text{on}\) and \(P_C = \sum_{m=1}^{M} \bar{e}_m P_C^\text{on}\), where \(P_A^\text{on}\) is the minimum power consumption of the PA and is equal to 28.75 mW, \(ν_A\) is the PA efficiency, which is considered as 0.5, and \(P_A^\text{off} = 4.375\text{mW}\). Furthermore, \(P_A^\text{on} = 150\text{mW}\) and \(P_A^\text{off} = 16\text{mW}\).

The power for signaling and control bits is \(P_H = \frac{E_{ref}}{\bar{e}_m} \sum_{k=1}^{K} R_k\), where \(E_{ref}^\text{H}\) is the energy intensity of the signaling and control given in [7], which equals to 0.35 J/Mbits and 0.06 J/Mbits for suburban environments with 0.75 UEs per km\(^2\) and 5.25 UEs per km\(^2\), respectively, and 1.99 J/Mbits and 0.3 J/Mbits for rural environments with 0.075 UEs per km\(^2\) and 0.525 UEs per km\(^2\), respectively.

\section{III. THE ADAPTIVE LSAS ALGORITHM}

\subsection{A. ALSAS DESIGN PRINCIPLE}

ALSAS aims to satisfy the UEs’ QoS requirements with the minimum transmit power by intelligently utilizing prior
knowledge of the requirements to adapt various system parameters. The adaptation process can be performed in two main stages as shown in Fig. 1. In the first stage, the main LSAS subsystems, which are responsible for signal transmission/reception as well as the UEs’ measurement, detect the main UEs’ properties, such as pathloss, interference information, etc., and pass them to the ADTx stage to determine the transmission scheduling accordingly.

In the second stage, the ADTx reduces the transmission time based on the UEs’ bit rate and latency requirements through a buffering system. Such requirements are considered because the UEs support various types of applications with different data rate requirements. For example, VoIP requires low bit rates in a continuous fashion, but ultra high definition (UHD) video streaming requires high data rates that is at least 40 Mbps, but it can be bursty if the user is streaming short videos [40]. Therefore, AAO can be used to reduce the number of active antennas without compromising the achievable bit rate by adjusting the transmission period according to a predefined criterion.

B. THE ADTx STAGE

The ADTx stage arranges the data into \( L \) scheduled transmissions (STs), and determines the precoding technique that minimizes the total required transmission time for each ST, subject to a given complexity constraint. Each ST contains data for one or more UE to be transmitted simultaneously, thus, resource allocation in the form of RB assignment is performed. For each ST, the allocation process specifies the number of required OFDM symbols \( \tilde{S} \in \mathbb{I}^{1 \times L} \), number of active antennas \( \mathcal{M} \in \mathbb{I}^{1 \times L} \), and the RB allocation matrix \( \mathcal{A} \in \mathbb{B}^{K \times N} \). The elements of \( \mathcal{A} \) and \( \tilde{S} \) are affected by the UEs’ requirements such as the data rates \( \tilde{R} \) and latencies \( \tilde{Y} \), and channel conditions including pathloss and time correlation.

The need for adaptive precoding can be justified by noting that CJ enables using all \( M \) antennas at the expense of significant multiple access interference. On the other hand, ZF eliminates the interference, but it limits the number of antennas that can be used due to the large delay associated with processing the precoding matrix [41]. In this paper, we assume that ZF can only use a subset of \( M_Z < M \) antennas. The choice of which precoding technique is used at the \( l \)th ST is influenced by \( M_Z \), the maximum number of multiplexed UEs when using CJ \( K_G \), or ZF \( K_Z \), as well as \( \psi \), \( \tilde{R} \) and \( \tilde{L} \), where \( M_Z < M \) and \( K_G > K_Z \). UEs that require transmission at the \( l \)th ST are indicated by the set \( \xi \in \mathbb{B}^{1 \times K} \), where \( \xi_k = 1 \) indicates that the \( k \)th UE requires an ST, otherwise \( \xi_k = 0 \). The ADTx stage during the \( l \)th ST can be best described using the following steps:

1) Set \( l = 1, \omega = 1. \)
2) Compute PO.
3) Select the desired Precoding scheme.
4) Allocate Resources.
5) If \( \sum_{k=1}^{K} R_{l,k} = 0 \) or \( \tilde{S}_{l+1} = 0 \), end.
6) Else, \( l = l + 1, L = l \), Go to step 2.

Steps 2, 3 and 4 are described below, and \( \omega \) is a weighting factor defined in Algorithm 3.

1) COMPUTING PO

The PO block provides all the required information for the precoder selection and resource allocation, which includes...
the SINR, number of RBs for each UE, achievable bit rate and required OFDM symbols. The achievable bit rate per OFDM symbol for CJ and ZF requires the SINRs of the CJ and ZF precoders at the lth ST for all users, which are given by \( y_{G,l} = \left[ y_{G,1}^{(1)}, y_{G,1}^{(2)}, \ldots, y_{G,1}^{(K,G,l)} \right] \) and \( y_{Z,l} = \left[ y_{Z,1}^{(1)}, y_{Z,1}^{(2)}, \ldots, y_{Z,1}^{(K,Z,l)} \right] \), respectively, where

\[
y_{G,k}^{(k)}(l) = y_{k|l}[M, K_{G,l}, a_{G,k}, \mu_G], \quad k = 1, \ldots, K_{G,l}
\]

\[
y_{Z,k}^{(k)}(l) = y_{k|l}[M_Z, K_{Z,l}, a_{Z_k}, \mu_Z], \quad k = 1, \ldots, K_{Z,l}
\]

and \( K_{G,l} \) and \( K_{Z,l} \) are the number of multiplexed UEs at the lth ST,

\[
K_{G,l} = \min \{ K_G, K_l \}
\]

and

\[
K_{Z,l} = \min \{ K_Z, K_l \}
\]

\( K_l \) is the total number of users whose data will be transmitted during the lth ST, \( \sum_{l=1}^{L} K_l = K \). Furthermore, the number of allocated RBs for the kth user during the lth ST using CJ precoding is given by

\[
N_{G,k,l} = \left[ \frac{K_{G,l}}{K_l} \right] B \xi_k, \quad k = 1, 2, \ldots, K_l
\]

subject to:

\[
\sum_{k=1}^{K_l} N_{G,k,l} = K_{G,l} B \xi_k
\]

\[
N_{G,k,l} \leq \left[ \frac{K_{G,l}}{K_l} \right] B + 1.
\]

Similarly, the number of RBs using ZF precoding is given by

\[
N_{Z,k} = \left( \frac{K_{Z,l}}{K_l} \right) B \xi_k
\]

subject to:

\[
\sum_{k=1}^{K} N_{Z,k} = K_{Z,l} B \xi_k
\]

\[
N_{Z,k} \leq \left[ \frac{K_{Z,l}}{K_l} \right] B + 1.
\]

Finally, using \( N_{G,k} \) and \( N_{Z,k} \), the calculated bit rate per OFDM symbol for \( k = 1, \ldots, K \) using CJ and ZF are given by

\[
R_{G,k} = R_{G}^{\max}(y_{G,k,l}, N_{G,k})
\]

\[
R_{Z,k} = R_{Z}^{\max}(y_{Z,k,l}, N_{Z,k})
\]

2) PRECODING SELECTION

In general, precoding selection leads to two possible outcomes, which are reducing the release time (RT) or increasing the spectral efficiency (SE). RT is defined as the time period taken for at least one UE to satisfy their bit rate requirement. Reducing RT benefits some UEs that significantly gain from a certain precoding choice, which in turn reduces the number of simultaneous UEs for the rest of the STs. However, this strategy may lead to blocking out the resource for the remaining UEs, especially at the cell edge. In order to overcome such a conflict, we propose the following:

- Precoding selection is based on minimizing the UEs’ outage probability.
- Trade-off between minimizing RT and maximizing SE where these two parameters are selected alternately only in the case when the RT of CJ and ZF are similar. If, however, RT for one of the precoding schemes is much smaller than the other, the precoding technique that minimizes RT is selected.

RT of CJ and ZF at the lth ST, is \( s_{G} = \min_0 \left( R_{t,o}/R_{G,0}^{\max} \right) \) and \( s_{Z} = \min_0 \left( R_{t,o}/R_{Z,0}^{\max} \right) \), respectively, where \( o = \{ k \mid k = 1, \ldots, K \land \xi_k = 1 \} \) and \( \land \) denotes the logical conjunction. In addition, the user that provides the minimum RT at the lth ST is given by \( \nu_{\min} = \arg \min \nu \). Hence we can define RT from CJ and ZF at the lth ST \( \nu = (s_{G}, s_{Z}) \) with its minimum and maximum values \( \nu_{\min} = \min \nu \) and \( \nu_{\max} = \max \nu \), respectively.

Minimizing RT takes the precoding technique that has the smallest value from set \( \nu \). Let the RT based precoding selection be denoted by \( P_I \in \{0, 1\} \), where \( P_I \) is the precoding indicator, \( P_I = 1 \) implies that ZF is used and \( P_I = 0 \) means that CJ is used. Thus, the precoding selection can be given by

\[
P_I = (\arg \min \{ s_{G}, s_{Z} \}) - 1.
\]

By noting that \( \arg \min \{ s_{G}, s_{Z} \} \in \{1, 2\} \), then \( P_I \in \{0, 1\} \). On the other hand, SE based precoding selection compares the bit rates of the precoding techniques to select the one with the largest value. If the outcome of SE based precoding selection is given by \( \Omega = 0 \), where \( \Omega = 1 \) corresponds to ZF and \( \Omega = 0 \) corresponds to CJ, then SE based precoding selection can be given by

\[
\Omega = U \left( \sum_{k=1}^{K} R_{G,k}^{\max}, \sum_{k=1}^{K} R_{Z,k}^{\max} \right).
\]
precoding technique that offers the minimum RT. Otherwise, the algorithm proceeds to check three and selects the precoding technique that maximizes SE, which will be performed using Algorithm 3. It is worth noting that Algorithm 3 guarantees that RT minimization is still considered when the overall outage probability for the users is minimized. This is achieved by using a weighting factor $\omega \in \{1, 2\}$, as shown in Algorithm 3.

The outcome of precoding selection at the $l$th ST will be defined as

$$P_l = \begin{cases} 1, & \text{ZF is selected} \\ 0, & \text{CJ is selected.} \end{cases}$$

(21)

Using the selected precoding technique, the global parameters for the $l$th ST, are given by $M_l = M$, $K_l = K_{G,l}$, $a_l = a_G$, $\mu_l = \mu_G$ and $N_l = N_G$ if $P_l = 0$; otherwise, $M_l = M_Z$, $K_l = K_{Z,l}$, $a_l = a_Z$, $\mu_l = \mu_Z$ and $N_l = N_Z$.

**Algorithm 3 Precoding Selection**

**Input:** $\mathbf{R}$, $\mathbf{R}_Z^{\max}$, $\mathbf{R}_G^{\max}$, $s_G$, $s_Z$, $S_l$, $\mathbf{v}$

1. if $s_G > S_l$ OR $s_Z > S_l$ keep as is
2. if $s_G = s_Z = 0$ then return $P_l = 0$
3. else
4. $P_l = P_1$, $S_l = s_{\min}$, $\omega = 1$
5. else
6. $P_l = P_3$, $S_l = s_{\max}$, $\omega = 1$
7. end
8. Output: $S_l$, $P_l$, $\omega$

**Algorithm 3 SE Based Precoder Selection**

**Input:** $\mathbf{R}_G^{\max}$, $\mathbf{R}_Z^{\max}$, $\mathbf{v}$, $\Omega$

1. $\lambda_Z = \sum_{i=1}^{L} U(\mathbf{R}_{Z,i}^{\max})$
2. $\lambda_G = \sum_{i=1}^{L} U(\mathbf{R}_{G,i}^{\max})$
3. if $\lambda_Z \neq \lambda_G$, then $P_l = U(\lambda_G, \lambda_Z)$
4. else, $P_l = \Omega$
5. end if
6. if $P_l \neq P_1$, then $\omega = 2$, $S_l = \max \left( \frac{\mathbf{v}P_l+1}{\omega}, \mathbf{v}P_l+1 \right)$
7. else, $\omega = 1$, $S_l = \mathbf{v}P_l+1$
8. end
9. Output: $P_l$, $S_l$, $\omega$

3) RESOURCE ALLOCATION

The final aspect in the ADTx stage is to utilize the information gathered to perform resource allocation which takes into account latency handling. After completing the precoding selection, resource re-allocation is needed if the latency constraints are not satisfied. A possible approach to overcome this is to increase the transmission time such that the latency constraint can be satisfied. However, this may reduce the system efficiency because it introduces transmission redundancy. Therefore, we propose Algorithm 4 to maintain redundancy within a pre-defined limit. Recalculating the number of RBs for the $l$th ST is given by $\tilde{N}_l$ subject to

$$\min \sum_{k=1}^{K} |R_{l,k} - \tilde{R}_k^{\max}|$$

$$\min \sum_{k=1}^{K} |\sum_{b=1}^{L} S_b \hat{z}_k - Y_k|.$$

In order to minimize the complexity of the proposed scheme, the maximum number of iterations is set to three, which is selected on the basis that the impact of resource re-allocation becomes negligible beyond three iterations.

Based on $\mathcal{N}_l$, UEs multiplexing and RBs allocation is given by

$$A_{l,k,n} = \begin{cases} 1, & \text{allocate } k\text{th UE at the } n\text{th RB} \\ 0, & \text{otherwise} \end{cases}$$

subject to:

$$\sum_{n=1}^{\mathcal{N}_l} A_{l,k,n} = N_{l,k}$$

$$\sum_{n=1}^{\mathcal{N}_l} A_{l,k,n} \leq K_{l}, \quad n = 1, \ldots, \mathcal{N}_l$$

where $\mathcal{A}_{\Sigma} \in \mathbb{R}^{1 \times \mathcal{N}_l}$, and $A_{\Sigma}^{k} = \sum_{n=1}^{\mathcal{N}_l} A_{l,k,n}$. $n = 1, \ldots, \mathcal{N}_l$. Using $S_l$ and $R_l$, the available symbols and remaining bit rates for the $(l + 1)$th ST are respectively given by $S_{l+1} = S_l - \sum_{b=1}^{L} S_b$ and $R_{l+1,k} = R_k - \sum_{b=1}^{L} R_{b,k}$.

**Algorithm 4 Latency Handling**

**Input:** $\mathcal{N}_l$, $R_{\max}$, $R_{\text{target},k}$, $Y$

1. $t = 1$
2. while $i \leq 3$ do
3. Adjust the number of RBs, $\tilde{N}_l$, based on the target, $R_{\text{target},k}$
4. Check the achievable bit rate, $\hat{R}_{\max}$, using $\tilde{N}_l$ set
5. If $\hat{R}_{\max} > R_{\max}$ for any $k = 1, \ldots, K$ then
6. $i = i+1$, reduce number of multiplexed UEs, $K_l$ and repeat line 3
7. else, $i = 4$, $N_{l,k} = \tilde{N}_{l,k}$ for $k = 1, \ldots, K$
8. end if
9. Output: $\tilde{N}_l$, $S_l$, and $K_l$

C. ANTENNA ARRAY OPTIMIZATION STAGE

The AAO stage attempts to reduce the number of active antennas required for each ST by adjusting the transmission period
based on \( R = S \log_2(1 + M \frac{\lambda}{d}) \). For example, assuming \( \frac{\lambda}{d} = 1 \) and \( M = 400 \), increasing the transmission time by 30% allows reducing the number of required antennas by 75%, and hence, provide further power saving. If the required number of antennas for the \( L \)th ST after the AAO is denoted as \( M_L \), the AAO attempts to find \( \min M_L \) subject to

\[
\sum_{l=1}^{L} \bar{S}_l \leq T_D \tag{24a}
\]

\[
\bar{R}_{l,k}^\text{max} S_l \leq \lambda_{l,k}, \quad k = 1, \ldots, K \tag{24b}
\]

\[
2K_l \leq \hat{M}_l \leq M \tag{24c}
\]

where \( \bar{S}_l \) is the number of OFDM symbols at the \( l \)th ST and \( \bar{R}_{l,k} \) is the bit rate for the \( k \)th UE during the \( l \)th ST after the AAO step. \( \bar{R}_{l,k} = \bar{S}_l C \left( \gamma \left( M_l, K_l, a_{l,k}, h_{l,k} \right), N \right) \), where

\[
C \left( \gamma (\cdot), N \right) = \frac{NQ}{T_f} \log_2 (1 + \gamma (\cdot)) \tag{25}
\]

IV. ADTx ANALYSIS

A. SINR ANALYSIS OF THE DTX TRANSMISSION

1) SINGLE CELL LSAS WITH \( K \) UES UNIFORMLY PLACED WITHIN A DISTANCE \( d_{\text{max}} \)

This scenario considers the case when the interference impact of LSAS to the surrounding cells is low. We assume that the pathloss is given by \( \alpha = 10^{(a_{\text{dB}}/10)} \), where \( a_{\text{dB}} = A + B \log_{10}(d) \), \( A \) is a constant that captures the effect of free space propagation at a reference distance \( d \), and \( B = 10\beta \), where \( \beta \) is the pathloss exponent. In the following analysis, the UE index is dropped, unless it is necessary to include it.

Lemma 1: The probability and cumulative density functions (PDF) and (CDF) of the UE’s SINR are respectively given by

\[
f_{\psi} (\psi) = \begin{cases} \frac{20}{\psi B_d^2} \left( \frac{\alpha}{\psi} \right)^{20}, & \psi \geq \psi_{\text{min}} \\ 0, & \text{otherwise} \end{cases} \tag{26}
\]

\[
F_{\psi} (\psi) = \begin{cases} \frac{1}{\alpha^2} \left( \left( \frac{\alpha}{\psi_{\text{min}}} \right)^{20} - \left( \frac{\psi}{\psi_{\text{min}}} \right)^{20} \right), & \psi \geq \psi_{\text{min}} \\ 0, & \text{Otherwise} \end{cases} \tag{27}
\]

where \( \sigma = P_T / (\eta 10^{A/10}) \), \( \eta \) is the noise level at the UE, and the SINR per antenna \( \psi = P_T / (\alpha \eta) \). The proof is given in Appendix II.

2) \( K \) UES TRANSMITTING USING \( L \) STS

Consider the case where \( K \) UEs are transmitting over \( L \) STs, and each ST consists of \( K_l \) UEs, \( \sum_{l=1}^{L} K_l = K \). The UEs are grouped based on their SINR values in ascending order, e.g. the UE with the lowest SINR is transmitted during the first ST, and the one with the highest SINR is transmitted in the \( L \)th ST. The UE whose SINR order is \( m \) during the \( m \)th ST is denoted as \( u_m \).

Lemma 2: The PDF of the SINR of the \( u_m \)th UE is given by

\[
f_{u_m} (\psi) = \frac{Kb}{r} \left( \frac{K - 1}{u_m - 1} \right) \sum_{i=0}^{\phi_m} \binom{\phi_m}{i} q^i \times \left( g \psi \frac{1}{u_m} \left( h_m + \frac{r}{k} \right) - \psi \frac{1}{u_m} \left( h_m + \frac{r}{k} \right) \right)^{h_m} \tag{28}
\]

where \( \phi_m = K - u_m \), \( b = \frac{\sigma}{\psi^2} / d_{\text{max}}^2 \), \( \lambda = r + K \), and \( g = v/b \).

**Proof:** Using order statistics, the SINR PDF of \( u_m \) UE is given by [42]

\[
f_{u_m} (\psi) = W_m f_\psi (\psi) \left( F_\psi (\psi) \right)^{u_m-1} (1 - F_\psi (\psi))^{K-u_m} \tag{29}
\]

where \( W_m = K (K-1)/(u_m-1) \). Expanding (29) gives

\[
f_{u_m} (\psi) = W_m \left( \frac{\sigma}{\psi} \right)^{\frac{1}{2}} \left( v - \sigma \left( \frac{\psi}{\sigma} \right)^{\frac{1}{2}} \right)^{u_m-1} \times \left( 1 - v + \sigma \left( \frac{\psi}{\sigma} \right)^{\frac{1}{2}} \right)^{K-u_m} \tag{30}
\]

where \( \sigma = 1/d_{\text{max}}^2 \). After some straightforward manipulations (30) can be written as

\[
f_{u_m} (\psi) = \frac{\Lambda_{u_m}}{\psi^{\frac{20}{20}}} \left( v - b \right)^{u_m-1} \left( 1 - v + b \right)^{K-u_m} \tag{31}
\]

where

\[
\Lambda_{u_m} = \frac{Kb}{r} \left( K - 1 \right) \tag{32}
\]

By defining \( \psi^{\frac{20}{20}} = x \), and noting that \( b > 0 \), \( \{u_m, \Lambda_{u_m}\} \) are integers greater than one, thus \( f_{u_m} (\psi) \) can be written as

\[
f_{u_m} (x^{\frac{20}{20}}) = \frac{\Lambda_{u_m}}{x^{\frac{20}{20}}} \left( v - b \right)^{u_m-1} \left( 1 - v + b \right) \left( x^{\frac{20}{20}} \right)^{u_m} \tag{33}
\]

which can be further simplified to

\[
f_{u_m} (x^{\frac{20}{20}}) = \frac{\Lambda_{u_m}}{x^r} \left( gx - 1 \right)^{b_m} (qx + 1)^{b_m} \tag{34}
\]

Using binomial expansion and substituting \( x^{\frac{20}{20}} \) gives the PDF of \( u_m \)th-ranked UE given in (28).

B. TRANSMISSION PERIOD OF ADTx

Under scenario 2 and assuming that all UEs require the same throughput, \( R \), the transmission length of the \( L \)th ST is affected by the lowest SINR for that particular group. This implies that

\[
S_l = \max_{u_m=1, \ldots, K_l} \frac{R}{C_l} \tag{35}
\]
TABLE 3. Simulation parameters for Fig. 2 and 3, [7].

| Parameter       | Suburban | Rural |
|-----------------|----------|-------|
| $\theta$        | 139.1    | 136.6 |
| $\beta$         | 3.76     | 3.86  |
| $d_{\min}$ (m)  | 35       | 35    |
| $d_{\max}$ (km) | 1        | 1.345 |
| $T_{fr}$ (ms)   | 45       | 45    |
| $B$             | 100      | 100   |

\[
C_I = \begin{cases} 
\frac{Q}{T_{fr}} \log_2 \left(1 + \frac{M}{K_l} \psi_{l,\text{um}} \left(\psi_{l,\text{um}} + 1\right)\right), & \text{CJ} \\
\frac{Q}{T_{fr}} \log_2 \left(1 + \frac{M_{Z} - K_i}{K_l} \psi_{l,\text{um}}\right), & \text{ZF}.
\end{cases} 
\]  
\( (36) \)

Hence, the average transmission period for scenario 2 is given by
\[
S = \sum_{l=1}^{L} S_l, 
\]
where
\[
S_l = \int_{-\infty}^{\infty} f_{\text{um}}(\psi) S_l(\psi) d\psi 
\]
\( (37) \)

Assuming that in each ST, adaptive precoding selection is applied based on the following condition:
\[
P_l = \begin{cases} 
1, & \psi_{l,\text{um}} \geq \psi_Z \\
0, & \psi_{l,\text{um}} < \psi_Z
\end{cases} 
\]
\( (39) \)

V. NUMERICAL RESULTS
A. IMPACT OF ADTx

The impact of discontinuous transmission on LSAS is evaluated in Figs 2 and 3 for the suburban and rural scenarios, respectively. LSAS is considered to have $M = 400$ antennas, where each antenna transmits 200 mW of power. The simulation parameters are given in Table 3. The channel attenuation factors include the building penetration [7], [43]. The analytical results are obtained using (40).

As can be noted from Figs 2 and 3, the analysis and simulation results match almost perfectly, which confirms the validity of the derived analysis. As can be seen from these figures, increasing the number of STs reduces the total number of occupied slots. The optimum number of STs vary depending on the number of active UEs. For example, the minimum number of transmission slots for $K = 5$ and $K = 9$ is achieved using 2 and 3 STs, respectively. This shows that the proposed ADTx can reduce the required transmission period and allows switching off some of the LSAS components.

B. EVALUATION PARAMETERS FOR MULTI CELL SCENARIOS

The performance evaluation considered in this subsection focuses on EE and the outage probability performance. The UEs outage represents the number of UEs that do not get
at least half of the required bit rate. The presented results correspond to suburban and rural scenarios, which make up most of the coverage area of a cellular network, and hence, the EE improvement is expected to be significant. A typical users’ density in suburban areas ranges between 0.75 and 5.25 UE per km², with a short inter-site-distances (ISD) of 1732.1 m. In rural areas, the users’ density ranges between 0.075 and 0.525 UE per km² with ISD of 2330 m [9].

A 49 cell wrap around is considered [36], the carrier frequency is 2 GHz, and the pathloss is given by $\alpha_k = 10^{\frac{-X}{10}} (d_k)^{\beta}$ where $X$ is log normal shadowing with zero mean and 6 dB standard deviation. The small scale fading is modeled as Rayleigh, and the noise floor at the MBS and UEs is set to -172 dBm/Hz and -165 dBm/Hz, respectively. Various cases of bit rate requirements are evaluated, which includes fixed and varying bit rate requirements. In the case of fixed bit rate requirements, the target bit rate is set to be in the range of 1 to 6 Mbps. On the other hand, in the case of varying bit rate requirements we assume a minimum bit rate of 100 kbps with the maximum bit rate between 1 to 15 Mbps. The bit rate is generated in a uniform distribution in logarithmic scale. For both cases, latency requirements between $10^{-3}$ and 5 duty cycles per frame is considered.

The performance of the proposed scheme is compared to the continuous transmission case described in Appendix A. The performance of the proposed ALSAS in a suburban scenario shown in Figs. 4, 5, and 6, where Fig. 4 and Fig. 5 considers the EE with fixed and variable bit rate requirements, while 6 presents the outage probability. As can be noted from Fig. 4, the proposed technique provides 125% EE improvement for a fixed bit rate requirement of 1 Mbps for the 5.25 users/km² case. When the user density decreases to 0.75 users/km², the EE improves substantially to 455%. As the bit rate requirement and number of UEs increase, the improvement decreases and becomes about 1% for the high density case, while it remains significant for the low density case with an improvement of about 52%. The EE improvement reduction at high data rate and user density is due to the fact that the MBS is mostly required to transmit over the entire transmission frame.

**FIGURE 4.** Energy productivity of suburban scenario with fixed rate requirement.

**FIGURE 5.** Energy productivity of suburban scenario with varying bit rate requirement.

**FIGURE 6.** Outage probability of suburban scenario with fixed rate requirement.
The EE for varying bit rate is shown in Fig. 5. As can be noted, the EE trends with variable bit rates are generally similar to the EE with fixed bit rates given in Fig. 4, except that the improvement is considerably higher, particularly for the low user density case. As can be noted from the figure, the proposed ALSAS, for the case of 0.75 UEs per km², provides improvement of 1124% and 118% for the maximum bit rate of 1 and 15 Mbps, respectively. For the high user density case, the EE improvement is 304% and 23% for 1 and 15 Mbps, respectively. Considering that most of the wireless mobile applications require bit rates less than 1.5 Mbps, the proposed technique will provide an energy saving of more than 1000% in this suburban scenario.

Fig. 6 shows the outage probability of the conventional and proposed adaptive LSAS techniques using fixed bit rate requirements. As the figure shows, ALSAS offers lower or equivalent outage probability for the considered bit rate range. Consequently, the EE improvement obtained using the ALSAS is achieved without compromising the outage probability.

Similar to the suburban scenario, the performance of the proposed ALSAS is depicted in Figs 7-10. The results are presented for fixed and variable bit rates and ISDs. The fixed ISD is used to evaluate the performance of the proposed technique with conventional configuration, while the varying ISD is used to evaluate the EE performance if the cellular operators of LSAS have ISD values that are from the one discussed in this paper. To evaluate the impact for various ISDs, the required fixed bit rate is set to 4 Mbps, and the varying bit rate ranges from 100 kbps to 10 Mbps. In both cases, the user density is considered as 0.525 UE/km².

Fig. 7 shows the EE for fixed ISD and fixed bit rates. Similar to the suburban scenario, the proposed technique improves the EE of the cellular network in rural scenarios without degrading the QoS. Fig. 7 shows that an EE improvement of more than 21% is guaranteed given that the bit rate requirement is less than 6 Mbps for all UEs. As can be seen from this figure, the proposed ALSAS provides a 61% EE for low loads such as 0.075 UE/km², even if the bit rate is 6 Mbps. When the load increases to 0.525 UE/km², the EE improvement surges to 196%.

Fig. 8 shows the EE for a fixed ISD and variable bit rates. As can be noted from this figure, the EE improvement for ALSAS at 0.75 UEs/km² is about 952% and 150% for a maximum bit rate of 1 and 15 Mbps, respectively. For the case of 0.525 UEs/km², the EE improvement is about 461% and 32% compared to conventional LSAS for a maximum bit rate of 1 and 15 Mbps, respectively. Considering that most applications require less than the simulated bit rate requirement, it can be concluded that a minimum of 150% EE improvement will be provided for future cellular networks.

The QoS in terms of outage probability is shown in Fig. 9 for the fixed rate and fixed ISD scenario. As depicted in this figure, the proposed ALSAS managed to offer an outage probability that is comparable to the one offered by the conventional LSAS for both low and high user densities.

The EE and outage probability for fixed and variable bit rates and ISDs are presented in Fig. 10 and Table 4, respectively. As can be noted from Fig. 10 and Table 4, conventional LSAS achieves a maximum EE at 2.6 and 3 km, while the proposed ALSAS improves EE as the ISD is reduced, without compromising the QoS. Such performance is obtained because reducing ISD decreases the total number of users and the overall pathloss between MBS and UEs. Consequently, the required transmission time and number of active antennas can be reduced. This further proves the capability of
the proposed technique to maximize the hibernation time of the MBS. The figure also shows that conventional transmission does not provide a significant EE improvement as ISD is varied. In practical scenarios, the bit rate requirement for certain UEs might be varying, thus, the proposed technique can provide up to 500% EE improvement over conventional LSAS. Although reducing ISD causes higher capital expenditure (CAPEX) for the cellular operator, it lowers the carbon footprint significantly, as well as improves the network performance due to shortening the distance.

### VI. CONCLUSION

This paper proposed an adaptive system to minimize the power consumption in LSAS systems. The adaptive system reduces the power consumption by combining ADTx with precoding selection and antenna optimization. The ADTx stage is responsible for dividing the transmission into L STs and selecting a suitable precoding technique per ST such that the total transmission time and latency are minimized. After ADTx, antenna array optimization is applied to further increase the duration of the hibernation period for a subset of antenna elements. The obtained analytical and simulation results showed that the proposed technique can provide EE improvements between 125% and 1124% in the suburban scenario, and between 196% and 952% in the rural scenario, without compromising QoS. It was also demonstrated that for rural environments with relatively smaller ISD values, the proposed technique can provide up to 500% power saving for fixed bit rate requirements.

### APPENDIXES

### APPENDIX I

**CONTINUOUS TRANSMISSION RESOURCE ALLOCATION**

For the sake of completeness, this section describes the continuous transmission scheme with bit rate and SINR adaptation. The adaptation is performed in two steps, which are the power control calculation and the RBs allocation. To maintain low complexity and delay, the continuous transmission scheme uses CJ precoding.

- **Power control calculation:**

  The power control is applied if the achievable bit rate for all \( K \) UEs is less than the target bit rate. With \( K \) active users, the initial number of RBs per UE is given by \( N_{\text{in}} = \frac{B}{K} \) and the initial throughput for all UEs is given by

  \[
  R_{\text{in}} = \{(\gamma_k, N) \mid T_0C(\gamma_k, N) \land k = 1, \ldots, K\} \tag{41}
  \]

  where \( C \) is given by (7). The transmit power coefficient is given by

  \[
  \varphi = \begin{cases} 
  1, & \text{if } \max_{k=1,\ldots,K} \Theta_k = \emptyset \\
  0, & \text{otherwise} 
  \end{cases} \tag{42}
  \]

  where

  \[
  \mathcal{D} = \left\{(R_{\text{in},k}, R_k) \mid R_{\text{in},k} > R_k \land k = 1, \ldots, K\right\} \tag{43}
  \]

  and \( \Theta \in \mathbb{R}^K \) is given by \( \Theta_k = \gamma_{\text{target},k} / \gamma_k, k = 1, \ldots, K \), and \( \gamma_{\text{target},k} \) is the target SINR, which is given by

  \[
  \gamma_{\text{target},k} = \frac{3k \left( \Phi_k + \varphi_k + \Phi_k \varphi_k + 1 \right)}{1 - 3k}, \quad k = 1, \ldots, K.
  \]
\[ \Phi_k = \frac{I_k}{\eta_k}, \]
\[ \gamma_k = \frac{K}{(M - K) \eta_k}. \]
\[ (44) \]

As can be noted from (42), the power control is applied only when the achievable bit rate is lower than the throughput requirement.

- **RBs allocation**

After applying power control, further resource adjustment can be applied by considering the RB allocation. UEs that achieve more than the required bit rate can apply further reduction in the allocated RBs. This reduces the number of multiplexed UEs in some part of the RBs, hence, increases the achievable bit rate for the rest of UEs. The number of RBs for continuous transmission, \( \mathbf{N} \in \mathbb{R}^K \) is given by

\[ \mathbf{N}_k = \frac{R_{\text{f}1}}{\log_2 (1 + \gamma_{\phi,k})}, \]

such that \( \sum_{k=1}^{K} \mathbf{N}_k \leq K \mathbf{N} \) and \( \mathbf{N}_k \leq \mathbf{N} \) for \( k = 1, \ldots, K \), where \( \gamma_{\phi,k} \in \mathbb{R}^K \) is the SINR coefficient with a power control coefficient \( \phi \). Using \( \mathbf{N} \), the RBs allocation \( \mathbf{A} \in \mathbb{R}^{K \times \mathbf{N}} \), is given by

\[ \mathbf{A}_{k,n} = \begin{cases} 1, & \text{allocate } k^{\text{th}} \text{ UE at the } n^{\text{th}} \text{ RB} \\ 0, & \text{otherwise} \end{cases} \]
\[ (45) \]

subject to:

\[ \sum_{n=1}^{\mathbf{N}} \mathbf{A}_{k,n} = \mathbf{N}_k \]
\[ (46a) \]

\[ \mathbf{A}_{n} \leq \mathbf{N}_1, \quad n = 1, \ldots, \mathbf{N} \]
\[ (46b) \]

where \( \mathbf{A}_\Sigma = \sum_{k} \mathbf{A}_{k,n} \).

**APPENDIX II**

**PROOF OF LEMMA 1**

Given that the PDF of the distance, \( f_D (d) \) between the MBS and EU is given by

\[ f_D (d) = \begin{cases} \frac{2d}{d^2_{\text{max}}}, & d \leq d_{\text{max}} \\ 0, & \text{otherwise}. \end{cases} \]
\[ (47) \]

By noting that \( a_{dB} = A + B \log (d) \), then \( a = 10\frac{A + B \log(d)}{10} \). Substituting \( a \) in \( \psi \) gives

\[ \psi = \frac{P_F}{\eta \times 10^{\frac{A + B \log(d)}{10}}} = \frac{P_F}{\eta \times 10^{\frac{A + B \log(d)}{10}}} = \frac{\sigma}{d^{\frac{B}{10}}}. \]
\[ (48) \]

Because generally \( B \) is greater than 10, then \( \psi \) is strictly decreasing by increasing \( d \). Therefore, the CDF of \( \psi \) is given by,

\[ F_\psi (\psi) = P (d \leq u) = 1 - F_D (u) \]
\[ (49) \]

where \( F_D (d) \) is the CDF of \( d \) and \( u = \left( \frac{a}{\sigma} \right)^{\frac{10}{B}} \). Consequently, \( f_\psi (\psi) \) can be obtained using the chain rule,

\[ f_\psi (\psi) = \frac{\partial}{\partial \psi} (1 - F_D (u)) \frac{\partial u}{\partial \psi} = -F_D (u) \frac{\partial u}{\partial \psi} \]
\[ (50) \]

By substituting \( \frac{\partial u}{\partial \psi} = -\frac{10}{B \sigma d_{\text{max}}} \psi_{\text{min}} \) and \( f_D (u) = \frac{2}{d_{\text{max}}} \left( \frac{\sigma}{\psi} \right)^{\frac{10}{B}} \) in (50) and simplifying the result, the PDF in (26) is obtained.

The CDF of the SINR is thus given by

\[ \int_{\psi_{\text{min}}}^{\psi} \frac{1}{20 \pi} du = B \left( \psi_{\text{min}}^{20 \pi} - \psi^{20 \pi} \right). \]
\[ (52) \]

By substituting (52) in (51) we obtain

\[ \int_{\psi_{\text{min}}}^{\psi} \frac{1}{20 \pi} du = B \left( \psi_{\text{min}}^{20 \pi} - \psi^{20 \pi} \right) \]
\[ (53) \]

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