Equitable 2-partitions of the Hamming graphs with the second eigenvalue

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Abstract

The eigenvalues of the Hamming graph $H(n,q)$ are known to be
\[ \lambda_i(n,q) = (q - 1)n - qi, \quad 0 \leq i \leq n. \]
The characterization of equitable 2-partitions of the Hamming graphs $H(n,q)$ with eigenvalue $\lambda_1(n,q)$ was obtained by Meyerowitz in \[15\]. We study the equitable 2-partitions of $H(n,q)$ with eigenvalue $\lambda_2(n,q)$. We show that these partitions are reduced to equitable 2-partitions of $H(3,q)$ with eigenvalue $\lambda_2(3,q)$ with exception of two constructions.

1 Introduction

An ordered $r$-partition $(C_1, \ldots, C_r)$ of the vertex set of a graph is called equitable if for any $i, j \in \{1, \ldots, r\}$ there is $S_{ij}$ such that any vertex of $C_i$ has exactly $S_{ij}$ neighbors in $C_j$. The elements of the partition are called cells. It is well-known (see e.g. \[6\]) that the eigenvalues of the matrix $S = (S_{ij})_{i,j \in \{1,\ldots,r\}}$, which is called the quotient matrix of the equitable partition, are necessarily eigenvalues of the adjacency matrix of the graph. An eigenvalue of an equitable partition is an eigenvalue of its quotient matrix.

Equitable partitions of various graphs are solutions for certain covering and optimization problems in coding, graph and design theories, projective geometry and etc. Such objects as 1-perfect codes, $(w - 1) - (n, w, \lambda)$-designs \[16\] and their q-analogues, including q-ary Steiner triple systems \[4\], spreads and Cameron-Liebler line classes in $PG(n,q)$ \[7\] can be defined in terms of equitable 2-partitions of Hamming, Johnson and Grassman graphs.

A subset of the vertex set of the graph is called a completely regular code if the distance partition with respect to the subset is equitable. Obviously any cell of an equitable 2-partition is a completely regular code. Completely regular

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codes in Hamming graphs include Preparata, some BCH and perfect codes. For a survey on completely regular codes in Hamming graphs we refer to [2].

In throughout the paper we index the eigenvalues of the graphs in the descending order, starting with zeroth. The eigenvalues of the Hamming graph \( H(n, q) \) are known to be \( \lambda_i(n, q) = (q - 1)n - qi, \) \( i \in \{0, \ldots, n\} \), see e.g. [5]. The characterization of all completely regular codes (in particular equitable 2-partitions) with the first eigenvalue was obtained for the Hamming and Johnson graphs by Meyerowitz in [15]. In the paper we consider the case when equitable 2-partitions of the Hamming graphs have the second eigenvalue.

In binary case the equitable 2-partitions were studied by Fon-Der-Flaas [9, 10]. Consider an equitable 2-partition of the Hamming graph \( H(n, 2) \) with an asymmetric quotient matrix and eigenvalues \( \lambda_i(n, 2) \) and \( n \). In [10] Fon-Der-Flaas showed that the number \( i \) of the minimum eigenvalue is not greater than \( 2n/3 \). Later in [11] he constructed 2-partitions which eigenvalue numbers attain the above bound. Krotov and Vorob’ev [12] studied the existence problem of such partitions and obtained a new necessary condition in terms of elements of quotient matrices. They also gave a characterization of such partitions in \( H(12, 2) \) with only one open case left. The database of equitable 2-partitions of binary Hamming graphs can be found in [13].

In case of arbitrary \( q \), the equitable 2-partitions of \( H(n, q) \) with the first eigenvalue could be described as those that can be obtained from 2-partitions of \( H(1, q) \) by adding \( n - 1 \) nonessential coordinate positions [15]. The case of linear equitable 2-partitions, i.e. such that the cells are cosets by a linear code of dimension \( n - 1 \) were characterized in [3] (see also [14]). Note that these partitions are related to Hamming codes and therefore generally have numbers of their eigenvalues exceeding 2.

In the paper [18] every eigenvector of \( H(n, q) \) with eigenvalue \( \lambda_i(n, q) \) was related to a set of eigenvectors of \( H(n - 1, q) \) with eigenvalue \( \lambda_{i-1}(n - 1, q) \) in a certain manner. Based on this connection the minimum weights of the eigenvectors of \( H(n, q) \) with eigenvalue \( \lambda_1(n, q) \) [18] and arbitrary eigenvalue [19] were found. In case of equitable partitions the approach of [18] relates every equitable 2-partition of \( H(n, q) \) with eigenvalue \( \lambda_i(n, q) \) to a set of eigenvectors of \( H(n - 1, q) \) with eigenvalue \( \lambda_{i-1}(n - 1, q) \) whose entries take values in the set \( \{0, 1, -1\} \). We characterize all eigenvectors with eigenvalue \( \lambda_1(n - 1, q) \) of the Hamming graph \( H(n - 1, q) \) taking values \( \{0, 1, -1\} \) in Section 4. The description of these vectors impose restrictions on the structure of the parent equitable 2-partition of \( H(n, q) \) with eigenvalue \( \lambda_2(n, q) \), which are shown to be equitable 2-partitions of \( H(3, q) \) in Section 4 up to several constructions that we give in Section 3. The basic theory concerning equitable partitions of the Hamming graphs and the details of the approach from [18] are given in Section 3.
2 Equitable partitions of the Hamming graphs

Let \(B_1, \ldots, B_n\) be finite sets, \(x\) and \(y\) be two tuples from the cartesian product \(B_1 \times \cdots \times B_n\). We say that \(x\) is \(s\)-adjacent to \(y\) if \(x\) and \(y\) differ only in \(s\)th coordinate. The vertex set of the Hamming graph \(H(n,q)\) is the cartesian \(n\)th power of a set \(A\) of size \(q\) and vertices \(x\) and \(y\) are adjacent if they are differ in exactly one coordinate position. Throughout the paper, \(\mathcal{N}\) denotes the set \(\{1, \ldots, n\}\).

A function on the vertex set of a graph is called a \(\lambda\)-eigenfunction if the vector of its values is an eigenvector of the adjacency matrix of the graph with eigenvalue \(\lambda\) or the all-zero vector.

Given a code (a set of vertices of a graph) \(C\), by \(\chi_C\) we denote the characteristic function of the code in the vertex set of the graph. Let \((u_1, \ldots, u_r)\) be an eigenvector of the quotient matrix \(S\) of an equitable partition \((C_1, \ldots, C_r)\) with eigenvalue \(\lambda\). Then it is easy to see that the function \(\sum_{i \in \{1, \ldots, r\}} u_i \chi_{C_i}\) is a \(\lambda\)-eigenfunction of the adjacency matrix of the graph, which is known as Lloyd’s theorem.

Theorem 1. \(\square\) An eigenvalue of an equitable partition of a graph is an eigenvalue of the graph.

Proposition 1. Let \((C, \overline{C})\) be an equitable 2-partition of a \(k\)-regular graph \(G\) with the quotient matrix \(S\). Then it is easy to see that the eigenvalues of \(S\) are \(k\) and \(S_{11} - S_{21}\). Moreover, the considerations prior to Lloyd’s theorem imply that \(\chi_C\) is the sum of a \(\lambda\)-eigenfunction (which is a constant uniquely defined by \(S\)) and a two-valued \((S_{11} - S_{21})\)-eigenfunction of \(G\) that is constant on \(C\) and \(\overline{C}\). It is well-known that if the characteristic function of a set in \(H(n,q)\) is orthogonal to any \(\lambda_j(n,q)\)-eigenfunction, for all \(j\) such that \(1 \leq j \leq t\), then the code is a \(t\)-orthogonal array (see e.g. \cite{6} [Theorem 4.4]). In particular, in the case when \((C, \overline{C})\) is an equitable 2-partition of \(H(n,q)\) with the second eigenvalue, the cell \(C\) is evenly distributed in the induced Hamming subgraphs \(H(n-1,q)\).

Let \((C, \overline{C})\) be an equitable 2-partition of a \(k\)-regular graph \(G\) with the quotient matrix \(S\). Then the following statements hold:

1. \(|C| = |V(G)| S_{21}/(S_{12} + S_{21})\) and the eigenvalues of the partition are \(k, S_{11} - S_{21}\).

2. If \(G = H(n,q)\) and \(S_{11} - S_{21} = \lambda_2(n,q)\), then for any \(i \in \mathcal{N}, \alpha \in A\) we have that

\[|\{x \in C : x_i = \alpha\}| = S_{21} q^{n-2}/2.\]

Proof. Double counting of edges in \(G\) between \(C\) and \(\overline{C}\) gives that \(|C| = S_{12} + S_{21}|V(G)|\). When \(G = H(n,q)\) and \(S_{11} - S_{21} = \lambda_2(n,q) = (q-1)n - 2q\) we have \(S_{11} + S_{12} = n(q-1)\) which gives \(|C| = S_{21} q^{n-1}/2\). Since \(C\) is a 1-orthogonal array, \(|\{x \in C : x_i = \alpha\}|\) is a constant number and we obtain the required.

Let \(f\) be a function defined on the vertices \(A^n\) of \(H(n,q)\), \(\alpha \in A, i \in \mathcal{N}\). Let \((f)_{i}^{\alpha}\) be the function such that \((f)_{i}^{\alpha}(y') = f(y)\), where \(y_i = \alpha\) and \(y'\) is obtained
Corollary 1. Let $H$ of equitable 2-partitions of the Hamming graph $H$ eigenfunction of (the characteristic function of a cell of the partition respectively).

3.1 Permutation switching construction

3 Constructions of equitable 2-partitions of $H(n,q)$ with the second eigenvalue

Let $f$ be a function defined on the cartesian product $B_1 \times \ldots \times B_n$. We say $i \in N$ is an essential coordinate of $f$, if there are $i$-adjacent vertices $x$ and $y$ such that $f(x) \neq f(y)$. A coordinate is essential for a 2-partition of $B_1 \times \ldots \times B_n$ if it is essential for the characteristic function of one of its cells. Given a function $f : A^n \rightarrow \mathbb{R}$, we define the function $f^+ : A^{n+1} \rightarrow \mathbb{R}$ by the rule $f^+(x_1, \ldots, x_n, x_{n+1}) = f(x_1, \ldots, x_n)$ for any $x_{n+1} \in A$.

The following result is a folklore, its proof is straightforward. Lemma in case of equitable 2-partitions of the Hamming graph $H(n,2)$ could be found in [9].

Lemma 2. [9] 1. Let $f$ be a real-valued function defined on $A^n$. Then $f$ is $\lambda_i(n,q)$-eigenfunction iff $f^+$ is $\lambda_i(n+1,q)$-eigenfunction.

2. Let $(C,\overline{C})$ be arbitrary 2-partition of $H(n,q)$, i be a nonessential coordinate of the partition. Then $(C,\overline{C})$ is equitable with eigenvalue $\lambda_i(n,q)$ iff the partition of $H(n-1,q)$ obtained by deleting $i$th position in all tuples of $C$ and $\overline{C}$ is equitable with eigenvalue $\lambda_i(n-1,q)$.

We say that a function defined on $B_1 \times \ldots \times B_n$ (2-partition of $B_1 \times \ldots \times B_n$ respectively) is reduced if all coordinates from $N$ are essential for the function (the characteristic function of a cell of the partition respectively).

3.1 Permutation switching construction

Here we present a construction of equitable 2-partitions of $H(n,q)$ from specific equitable 2-partitions of $H(2,q)$ by switchings of coordinate positions.

Recall that the Cartesian product $G \square H$ of graphs $G$ and $H$ is a graph with the vertex set $V(G) \times V(H)$; and any two vertices $(u, u')$ and $(v, v')$ are adjacent
if and only if either \( u = v \) and \( u' \) is adjacent to \( v' \) in \( H \), or \( u' = v' \) and \( u \) is adjacent to \( v \) in \( G \).

Firstly, we describe all equitable 2-partitions of \( H(1,q)\Box H(1,q') \) with the smallest eigenvalue.

**Proposition 2.** A 2-partition \((C,\overline{C})\) of \( H(1,q)\Box H(1,q') \) is equitable with the quotient matrix \( S \) and eigenvalue \(-2\) if and only if \(|C \cap K|/|K| = S_{21}/(q+q')\) for any maximal clique \( K \) of \( H(1,q)\Box H(1,q') \).

**Proof.** Let \( C \) be a subset of \( H(1,q)\Box H(1,q') \) such that any maximal clique \( K \) contains exactly \(|K|\) vertices of \( C \). Then the partition \((C,\overline{C})\) is equitable with the quotient matrix \[
\begin{pmatrix}
a(q+q') - 2 & (1-a)(q+q') \\
(a(q+q') - 2) & (1-a)(q+q') - 2
\end{pmatrix}
\]
and has eigenvalue \(-2\).

If \( K \) is a maximal clique in \( H(1,q)\Box H(1,q') \), then it is easy to see that the partition \((K,\overline{K})\) is an equitable 2-partition of \( H(1,q)\Box H(1,q') \) with the quotient matrix \[
\begin{pmatrix}
|K| - 1 & q + q' - 1 - |K| \\
1 & q + q' - 3
\end{pmatrix}
\]
that has eigenvalue \(|K| - 2\). Let \((C,\overline{C})\) be an equitable partition of \( H(1,q)\Box H(1,q') \) with the eigenvalue \(-2\). Then \( \chi_K = f_0 + f_1 \), \( \chi_C = h_0 + f_2 \), where \( f_0 \) and \( h_0 \) are constant functions that are uniquely determined by the quotient matrices of the partitions, \( f_1 \) is a \(|K| - 2\)-eigenfunction and \( f_2 \) is a \((-2\))-eigenfunction. Since eigenvectors of the adjacency matrix with different eigenvalues are orthogonal, we have that:

\[
|C \cap K| = \sum_{v \in V(H(1,q)\Box H(1,q'))} \chi_K(v) \chi_C(v) = \sum_{v \in V(H(1,q)\Box H(1,q'))} h_0(v) f_0(v),
\]
which does not depend on \( K \), but only on \(|K|\), since the value of \( f_0 \) is uniquely determined by the quotient matrix of the partition \((K,\overline{K})\). Since the graph \( H(1,q)\Box H(1,q') \) could be parted into maximal cliques of the size \(|K|\), the number \(|C \cap K|/|K|\) is equal to \(|C|/qq'\). Finally, from Proposition 2 we have \(|C| = S_{21}qq'/(S_{12} + S_{21})\), which taking into account that \( S_{12} + S_{21} = q + q' - 2 - S_{11} + S_{21} = q + q' \) implies that the percentage of \( C \) in \( K \) is \( S_{21}/(q+q') \).

\[\square\]

**Construction A.** Let \( A_1, \ldots, A_{n-1} \) be a partition of \( A \). The vertices of \( H(2,q) \) are parted into \( A_i \times A, i = 1, \ldots, n-1 \). Consider a 2-partition \((C,\overline{C})\) of \( H(2,q) \) such that \(|C \cap K|/|K|\) is the same for any maximal clique \( K \) of any subgraph induced by \( A_i \times A, i = 1, \ldots, n-1 \). By Proposition 2 the restriction of \((C,\overline{C})\) to any of these subgraphs is equitable. Moreover, \((C,\overline{C})\) is an equitable 2-partition of \( H(2,q) \) with the eigenvalue \( \lambda_2(2,q) = -2 \) as any maximal clique of \( H(2,q) \) is a union of maximal cliques of the subgraphs.

Define \( C^+ \) to be obtained from \( C \) by adding \( n-2 \) nonessential coordinate positions, i.e. \( C^+ = \{(x_1, x_2, y_3, \ldots, y_n) : (x_1, x_2) \in C, y_i \in A, i \in \{3, \ldots, n\}\} \).

From Lemma 2, the partition \((C^+,\overline{C^+})\) of \( H(n,q) \) is equitable with eigenvalue \( \lambda_2(n,q) \). The construction allows switchings of coordinates to be applied. Given
a permutation $\pi$ of the coordinates from $\mathcal{N}$ and $x \in \mathcal{A}^n$ by $\pi(x)$ denote the tuple $(x_{\pi(1)}, \ldots, x_{\pi(n)})$, for a subset $M$ of $\mathcal{A}^n$ denote by $\pi(M)$ the set $\{\pi(x) : x \in M\}$.
Let $\pi_i$ be the transposition $(2, i+1)$, and $\pi_1$ be the identity permutation. Define $C^+_\pi$ to be
\[
\bigcup_{i=1,\ldots,n-1} \pi_i(A_i \times \mathcal{A}^{n-1} \cap C^+).
\]

Note that the essential coordinates of $(C_\pi, \overline{C_\pi})$ are $\mathcal{N}$, while those of $(C, \overline{C})$ are 1 and 2. The construction above is somewhat similar to the construction of nonsystematic binary 1-perfect codes by consecutive switchings of $i$-components [1].

**Theorem 2.** The partition $(C^+_\pi, \overline{C^+_\pi})$ is an equitable 2-partition of $H(n, q)$ with the same quotient matrix as that of $(C^+, \overline{C^+})$ and eigenvalue $\lambda_2(n, q)$.

**Proof.** Let $S$ be the quotient matrix of the partition $(C^+, \overline{C^+})$. Define $(x)_\pi$ to be $\pi_i(x)$, for all $i \in \mathcal{N}\setminus n$, $x \in A_i \times \mathcal{A}^{n-1}$. The mapping $(\cdot)_\pi$ permutes the second and $(i+1)$th coordinate positions for vertices in $A_i \times \mathcal{A}^{n-1}$, so its restriction to the subgraph induced by $A_i \times \mathcal{A}^{n-1}$ is an automorphism of the subgraph. The permutation $\pi_i$ fixes 1, so the restriction of $(\cdot)_\pi$ to the subgraph induced by $A_i \times \mathcal{A}^{n-1}$ acts on the set of maximal cliques consisting of 1-adjacent vertices of the subgraph. By the choice of the initial partition $(C, \overline{C})$, the percentages of $C^+$ in all maximal cliques consisting of pairwise 1-adjacent vertices of the subgraph are the same and are equal to $S_{21}/(S_{12} + S_{21})$. Since $(C^+)\pi$ is $C_\pi$ and $(\cdot)_\pi$ permutes these cliques, the percentage of $C_\pi$ in these cliques is $S_{21}/(S_{12} + S_{21})$.

For a vertex $x \in C^+$, $x_1 \in A_i$ we see that $(\cdot)_\pi$ maps the set of all $j$-neighbors from $C^+$, for all $j \geq 2$, to the set of all $j$-neighbors of $(x)_\pi$ from $C_\pi$ because they are all in $A_i \times \mathcal{A}^{n-1}$. A maximum clique of $H(n, q)$ formed by pairwise 1-adjacent vertices is the union of maximal cliques of the subgraphs induced by $A_i \times \mathcal{A}^{n-1}$, $i \in \mathcal{N}\setminus n$. Since the percentages of $C$ and $C_\pi$ in the cliques in subgraphs are the same, we conclude that the numbers of vertices that are 1-adjacent to $x$ from $C^+$ and that of $(x)_\pi$ from $C_\pi$ coincide, so $(x)_\pi$ is adjacent to $S_{11}$ vertices of $C_\pi$. The proof in case when $x \in \overline{C^+}$ is analogous. 

\[\square\]

### 3.2 Alphabet liftings of two induced cycles in $H(4,2)$

Consider the construction of equitable partitions using alphabet liftings by Vorobiev in [20], which resembles the Zinoviev construction for perfect codes [21] (see also [17]). Let $A_0, \ldots, A_{q'-1}$ be the sets of the same size that partition $\mathcal{A}$. Here we identify the alphabet set $\mathcal{A}'$ of $H(n, q')$ with $\{0, \ldots, q'-1\}$.

Given a partition $(C_1, \ldots, C_r)$ of $H(n, q')$ define $D_i$, $1 \leq i \leq r$ to be
\[
\bigcup_{(x_1, \ldots, x_n) \in C_i} A_{x_1} \times \ldots \times A_{x_n}.
\]
Theorem 3. [20] Let $(C_1, \ldots, C_r)$ be an equitable partition of $H(n, q')$ with the set of eigenvalues $\{\lambda_i(n, q') : i \in I\}$. Then the partition $(D_1, \ldots, D_r)$ is an equitable partition of $H(n, q)$ with the set of eigenvalues $\{\lambda_i(n, q) : i \in I\}$.

Let $C = \{(0001), (0011), (0010), (0110), (1110), (1100), (1101), (1001)\}$. The complement $\overline{C}$ of $C$ is $\{(0000), (0100), (0101), (0111), (1111), (1011), (1010), (1000)\}$. We see that both $C$ and $\overline{C}$ are induced cycles in $H(4, 2)$, so any vertex of $C$ is adjacent to exactly two vertices of $\overline{C}$ and the partition $(C, \overline{C})$ is equitable with the quotient matrix $\begin{pmatrix} 2 & 2 \\ 2 & 2 \end{pmatrix}$. Moreover, all 4 coordinates of the partition are essential and $\lambda_2(4, 2) = 0$ is its eigenvalue.

Corollary 2. (construction B) Let $(C, \overline{C})$ be a partition of $H(4, 2)$ into two induced cycles of length 8. Then $\lambda_2(4, q) = 0$ is its eigenvalue.

4 Main results

Here we obtain the description of $\lambda_1(n, q)$-eigenfunctions with the values in $\{0, 1, -1\}$ and utilize it for the reconstruction of equitable 2-partitions of $H(n, q)$ with eigenvalue $\lambda_2(n, q)$.

4.1 $\lambda_1(n, q)$-eigenfunctions with the values in $\{0, 1, -1\}$

Consider two functions on the vertices of Hamming graphs. Given $i \in \mathcal{N}$, $A, B \subset A$, $A \cap B = \emptyset$, $A \cup B \neq \emptyset$ we say that $f$ defined on $A^n$ is the $(A, B, i)$-quasi string if

$$f(x) = \begin{cases} 
1, & x_i \in A \\
-1, & x_i \in B \\
0, & \text{otherwise}
\end{cases}$$

Given $i, j \in \mathcal{N}$, $i \neq j$, and two nonempty subsets $A, B$ of $A$ we say that $f$ defined on $A^n$ is the $(A, B, i, j)$-quasi cross if

$$f(x) = \begin{cases} 
1, & x_i \in A, x_j \notin B \\
-1, & x_i \notin A, x_j \in B \\
0, & \text{otherwise}
\end{cases}$$

Lemma 3. Let $f$ be a function from $A^n$ to $\{-1, 0, 1\}$. The function $f$ belongs to $U_0(n, q) \oplus U_1(n, q)$ iff it is the $(A, B, i, j)$-quasi cross or the $(A, B, i)$-quasi string or a constant.

Proof. Let us prove this lemma by induction on $n$. If $n = 1$, then it is easy to see that $f$ is either a constant or the $(A, B, i)$-quasi string.
Let us prove the induction step. Fix $\beta \in \mathcal{A}$. Lemma 1 implies that $f_n^\alpha - f_n^\beta \in U_0(n, q)$ for any $\alpha \in \mathcal{A} \setminus \beta$. Hence $f_n^\alpha - f_n^\beta$ is a constant, which we denote by $c_\alpha$, so we have that

$$f_n^\alpha \equiv f_n^\beta + c_\alpha. \quad (1)$$

Now, we consider three cases. Given a function $g$ by $E(g)$ we denote the set of its values.

**Case 1.** In this case we suppose that $\{-1, 1\} \subseteq E(f_n^\beta)$. Using the equality (1), we see that $E(f_n^\alpha)$ contains numbers $c_\alpha - 1$ and $c_\alpha + 1$. Since $E(f) \subseteq \{-1, 0, 1\}$, we have $c_\alpha = 0$ for any $\alpha \in \mathcal{A} \setminus \{\beta\}$. So $f_n^\alpha \equiv f_n^\beta$ for any $\alpha \in \mathcal{A} \setminus \{\beta\}$. By Lemma 1 we obtain $f_n^\beta \in U_0(n-1, q) \oplus U_1(n-1, q)$. Then using the induction assumption for $f_n^\beta$, we finish the proof in this case.

**Case 2.** Let $E(f_n^\beta)$ be $\{0, 1\}$ or $\{-1, 0\}$. Without loss of generality, we assume that $E(f_n^\beta) = \{0, 1\}$. Using (1) and the fact that $E(f) \subseteq \{-1, 0, 1\}$, we obtain that $c_\alpha \in \{-1, 0\}$ for any $\alpha \in \mathcal{A} \setminus \{\beta\}$. Denote $C = \{\alpha \in \mathcal{A} \setminus \{\beta\} : c_\alpha = -1\}$.

By Lemma 1 we obtain $f_n^\beta \in U_0(n-1, q) \oplus U_1(n-1, q)$. Using the induction assumption for $f_n^\beta$, we obtain that $f_n^\beta$ is the $(A, B, i)$-quasi string. Moreover, $B = \emptyset$ because $E(f_n^\beta) = \{0, 1\}$. Then $f$ is the $(A, C, i, n)$-quasi cross due to the definition of the cross.

**Case 3.** Let $f_n^\beta$ be a constant. Using (1), we obtain that $f_n^\beta$ is a constant for any $\alpha \in \mathcal{A} \setminus \{\beta\}$. Then either $f$ is a constant, or $f$ is the $(A, B, n)$-quasi string.

Let us now prove the sufficiency part. Obviously, the $(A, B, 1)$-quasi string belongs to $U_0(1, q) \oplus U_1(1, q)$ because the latter is the space of all real-valued functions on $H(1, q)$. The $(A, B, i)$-quasi string in $H(n, q)$ is obtained from the $(A, B, 1)$-quasi string in $H(1, q)$ by adding $n-1$ nonessential coordinate positions, so by Lemma 2 it is in $U_0(n, q) \oplus U_1(n, q)$.

On the other hand, it is easy to see that the $(A, B, i, j)$-quasi cross is the sum of the $(A, \overline{\mathcal{A}}, i)$-quasi string and the $(\overline{B}, B, j)$-quasi string that belong to $U_0(n, q) \oplus U_1(n, q)$, so the $(A, B, i, j)$-quasi cross is also in $U_0(n, q) \oplus U_1(n, q)$. \(\square\)

It is easy to see that the $(A, B, i)$-quasi string is in $U_1(n, q)$ iff $|A| = |B|$. If $x$ is a zero of the $(A, B, i, j)$-quasi cross then $x$ is $i$-adjacent to $|A|$ vertices $y$ such that $f(y) = 1$ and $|A|$ vertices $y$ such that $f(y) = -1$. If $x$ is such that $f(x)$ is positive (or negative) then $x$ is $i$-adjacent or $j$-adjacent to $|A| + q - |B| - 2$ vertices $y$ that have positive (or negative) value $|B| + q - |A| - 2$. Again, taking into account Lemma 2 the $(A, B, i, j)$-quasi cross is $\lambda_1(n, q)$-eigenfunction iff $|A| = |B|$. When $|A| = |B|$ the $(A, B, i)$-quasi string is called the $(A, B, i)$-string and the $(A, B, i, j)$-quasi cross is called the $(A, B, i, j)$-cross. The previous lemma implies the following.

**Lemma 4.** Let $f$ be an arbitrary function from $\mathcal{A}^n$ to $\{0, -1, 1\}$. The function $f$ is $\lambda_1(n, q)$-eigenfunction iff it is the $(A, B, i, j)$-cross or the $(A, B, i)$-string or the all-zero function.
4.2 Equitable 2-partitions of $H(n,q)$ with eigenvalue $\lambda_2(n,q)$

**Theorem 4.** Let $(C, \overline{C})$ be an equitable partition of $H(n,q)$ with the quotient matrix $S$ that has eigenvalue $\lambda_2(n,q)$. If $(\chi C)^n_k - (\chi C)^n_k$ is the $(A,B,i,j)$-cross and there is $s \in \mathcal{N} \setminus \{i,j,k\}$ that is essential for $(\chi C)^n_k$ or $(\chi C)^n_k'$ then $(C, \overline{C})$ is obtained by Construction B.

**Proof.** Without restriction of generality, suppose that $k = n$. Moreover, we have $|A| = |B|$, so up to an isomorphism of the Hamming graph we assume that $A = B$ and $(\chi C)^n_k - (\chi C)^n_k$ is the $(A,A,i,j)$-cross. Then we have the following properties.

\[ \{x : x_n = \alpha', x_i \in A, x_j \in \overline{A}\}, \{x : x_n = \alpha, x_i \in \overline{A}, x_j \in A\} \subset C, \quad (2) \]

\[ \{x : x_n = \alpha, x_i \in A, x_j \in \overline{A}\}, \{x : x_n = \alpha', x_i \in \overline{A}, x_j \in A\} \subset \overline{C}. \quad (3) \]

Let $x$ and $y$ be $n$-adjacent vertices such that $x_n = \alpha, y_n = \alpha'$, $x_i = y_i \in A, x_j = y_j \in A$. Then $x$ and $y$ are both in $C$ or not. \quad (4)

Let $x$ and $y$ be $n$-adjacent vertices such that $x_n = \alpha, y_n = \alpha'$, $x_i = y_i \in \overline{A}, x_j = y_j \in \overline{A}$. Then $x$ and $y$ are both in $C$ or not. \quad (5)

By conditions of the theorem there is a pair of $s$-adjacent vertices $x^{0000} \in C$ and $x^{0001} \in \overline{C}$ such that their $n$th positions are both $\alpha'$ or $\alpha$ and their $i$th and $j$th positions are in $A$ or not simultaneously by properties (2), (3). By properties (2) and (3) the essential coordinate positions for $(\chi C)^n_k$ and $(\chi C)^n_k'$ coincide, so we can assume that $x^{0000}_n = x^{0001}_n = \alpha'$.

Moreover, we assume that both $x^{0000}_i = x^{0001}_i = \beta', x^{0000}_j = x^{0001}_j = \gamma'$ are in $A$ (the case when they are in $A$ is proven analogously). W.l.o.g. up to an automorphism of $H(n,q)$ we assume that $\alpha$ and $x^{0001}_s = \delta \in A, \alpha'$ and $x^{0001}_s = \delta' \in \overline{A}$. By $\beta$ and $\gamma$ denote any two elements of $A$, so we have that $\alpha, \beta, \gamma, \delta \in A, \alpha', \beta', \gamma', \delta' \in \overline{A}$.

Denote by $x^{0101} a_1 a_2 a_3 a_4$ the vertex obtained from $x^{0001}$ by changing its $n$th position from $\alpha'$ to $\alpha$ iff $a_1 = 1$, its $i$th position from $\beta'$ to $\beta$ iff $a_2 = 1$, its $j$th position from $\gamma'$ to $\gamma$ iff $a_3 = 1$, its $s$th position from $\delta'$ to $\delta$ iff $a_4 = 1$. The graph spanned by $\{x^a : a \in \{0,1\}\}^4$ is $H(4,2)$. We now show that the partition $(C, \overline{C})$ could be reconstructed on the subgraph.

**Lemma 5.** $x^a \in C$ iff $a \in \{0000, 0100, 0101, 1000, 1010, 1011, 0111, 1111\}$

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Proof. Note that by (4) and (5) for any \( a \in \{0, 1\}^3 \) the vertex \( x^{0a} \) is in \( C \) iff \( x^{1a} \) is in \( C \) when \( x^{0a}, x^{0a} \in A \) or \( x^{0a}, x^{0a} \in \overline{A} \). Using this and taking into account (2) and (3) we have that

\[
x^0 \in C \text{ if } a \in \{0000, 0100, 0110, 1000, 1010, 1011\}, \tag{6}
\]

\[
x^0 \in \overline{C} \text{ if } a \in \{0001, 0010, 0011, 1001, 1100, 1101\}. \tag{7}
\]

We now show that \( x^{0110} \) is in \( \overline{C} \) and \( x^{0111} \) is in \( C \). Denote by \( x^{a_1a_2a_3} \) the vertex obtained from \( x^{a_1a_2a_3a_4} \) by deleting its \( j \)th coordinate position.

Consider the values of the function \( (\chi_C)^j_j - (\chi_C)^j_j \) on \( \{x^{a_1a_2a_3} : a_l \in \{0, 1\}, l = 1, 2, 3\}, \gamma = x^{0000}, \gamma = x^{0010} \). From (6) and (7) we have that

\[
((\chi_C)^j_j - (\chi_C)^j_j)(x^{000}) = \chi_C(x^{000}) - \chi_C(x^{0010}) = 1,
\]

\[
((\chi_C)^j_j - (\chi_C)^j_j)(x^{001}) = \chi_C(x^{001}) - \chi_C(x^{0011}) = 0,
\]

\[
((\chi_C)^j_j - (\chi_C)^j_j)(x^{100}) = \chi_C(x^{100}) - \chi_C(x^{1010}) = 0.
\]

If \( x^{0110} \) is in \( C \) then

\[
((\chi_C)^j_j - (\chi_C)^j_j)(x^{010}) = \chi_C(x^{010}) - \chi_C(x^{0110}) = 0.
\]

If \( x^{0111} \) is not in \( C \) then

\[
((\chi_C)^j_j - (\chi_C)^j_j)(x^{011}) = \chi_C(x^{011}) - \chi_C(x^{0111}) = 1.
\]

Therefore, if \( x^{0110} \) is in \( C \) or \( x^{0111} \) is in \( \overline{C} \) then \( (\chi_C)^j_j - (\chi_C)^j_j \) has at least three essential coordinate positions. However, by Corollary \( \Pi (\chi_C)^j_j - (\chi_C)^j_j \) is a \( \lambda_1(n-1, q) \)-eigenfunction of \( H(n-1, q) \), so it has not more then two essential coordinates according to the characterization in Lemma 4. Therefore we have that \( x^{0110} \in \overline{C} \) and \( x^{0111} \in C \) and then \( x^{1110} \in \overline{C} \) and \( x^{1111} \in C \) since changing \( n \)th position to \( 0 \) preserves the property of being in \( C \) in this case by properties (4) and (5).

\[\square\]

**Lemma 6.** We have that \( S_{12} = S_{21} = q, |A| = q/2 \) and the following holds up to an isomorphism of \( H(n, q) \).

A vertex \( x \) such that \( x_i, x_j \in A \) is in \( C \) iff \( x_s \in \overline{A} \). \( \tag{8} \)

A vertex \( x \) such that \( x_i, x_j \in \overline{A} \) is in \( C \) iff \( x_s \in A \). \( \tag{9} \)

Proof. Let \( A' \) be \( \{x' : x' \in \overline{C} \text{ is s-adjacent to } x^{0000}\} \). Lemma 5 holds for any vertex \( x^{0001} \in \overline{C} \) that is s-adjacent to \( x \), so we have that

\[
A \setminus A' = \{x_s : x \in \overline{C} \text{ is s-adjacent to } x^{0111}\}. \tag{10}
\]
We now evaluate the numbers of the neighbors of $x^{0000}, x^{0111}$ from $\overline{C}$. By the definition of equitable partition, it is $S_{12}$, where $(S_{ij})_{i,j=1,2}$ is the quotient matrix of the partition $(C, \overline{C})$. Consider the vertices obtained from $x^{0000} \in C$ by changing its $j$th symbol to an element from $A$. By property (3) these vertices are not in $C$, so there are at least $|A| + |A'|$ vertices from $\overline{C}$, that are $s$- or $j$-adjacent to $x^{0000}$. By property (8) the vertices obtained from $x^{0111} \in C$ by changing its $i$th symbol to an element of $A$ are in $\overline{C}$. The above combined with (10) gives that

$$S_{12} \geq \max\{|A| + |A'|, 2q - |A'| - |A|\}. \quad (11)$$

Now count the neighbors of $x^{0001}$ and $x^{0110}$ in $C$. For $x^{0001}$ there are $|A|$ $i$-neighbors and $q - |A'|$ $s$-neighbors from $C$. For $x^{0110}$ there are $q - |A|$ $j$-neighbors and $|A'|$ $s$-neighbors from $C$. Then we have that

$$S_{21} \geq \max\{q + |A| - |A'|, q - |A| + |A'|\}. \quad (12)$$

Since $\lambda_2(n, q) = n(q - 1) - 2q = S_{12} - S_{21}$ is an eigenvalue of the equitable partition with the quotient matrix $S$ we have that $S_{12} + S_{21} = 2q$, which combined with the inequalities (11) and (12) implies that $|A'| = |A| = q/2$ (in below up to an automorphism of $H(n, q)$ we assume that $A' = A$ and $S_{12} = S_{21} = q$). Moreover, the bound (11) is attained. This implies that the neighbors of $x^{0000}$ from $\overline{C}$ as exactly those $j$- and $s$-neighbors of $x^{0001}$ counted while obtaining (11). So all $l$-adjacent vertices of $x^{0001}$, $l \in N \setminus \{j, s\}$ are in $C$. Analogously, the vertices from $\overline{C}$ adjacent to the vertex $x^{0001}$ are exactly those $i$- and $s$-adjacent vertices to $x^{0001}$ counted while obtaining bound (12). We conclude that all vertices that are $l$-adjacent vertices to $x^{0001}$, $l \in N \setminus \{i, s\}$ are in $\overline{C}$. The consideration above holds for any two $s$-adjacent vertices $x^{0000} \in C$ and $x^{0001} \in C$ and we know their neighbors from $\overline{C}$ and $C$ respectively, so (8) and (9) hold.

By the previous Lemma, (2) and (1) the partition $(C, \overline{C})$ is reconstructed on the following set $\{x \in A^n : x_n = \alpha, \alpha'\}$. We now show that for any $\beta \in A$ the vertices of $\{x \in A^n : x_n = \beta\} \cap C$ are obtained by "copying" their $n$-neighbors either from $\{x : x_n = \alpha\} \cap C$ or from $\{x : x_n = \alpha'\} \cap C$.

Consider $f = (\chi_C)^{\alpha} - (\chi_C)^{\alpha'}$, $\beta \in A \setminus \{\alpha, \alpha'\}$. Then we have that

$$f(y) = 0 \text{ if } y_i, y_j \in A \text{ or } y_i, y_j \in \overline{A}. \quad (13)$$

We now show that $f$ (and therefore $(\chi_C)^{\alpha'}$) could be reconstructed in only two ways.

From (2), we have that $f(y) \leq 0$ if $y_i \in A$, $y_j \in \overline{A}$. We show that $f(y)$ is a constant for all $y \in A^{n-1}$, $y_i \in A$, $y_j \in \overline{A}$. Let $y$ be such that $f(y) = -1$, $y_i \in A$, $y_j \in \overline{A}$. Then from (15) $i$ and $j$ are essential coordinates for $f$ and by Lemma 3 we see that $f$ is the $(A', B', i, j)$-cross. Moreover, if $z$ is $i$- or $j$-adjacent to $y$, $f(z) = 0$, $z_i \in A$, $z_j \in \overline{A}$ then by (15) a nonzero $y$ of the function $f$ is adjacent to at least $q + 1$ zeros of $f$, which contradicts the definition of the
\((A', B', i, j)\)-cross. Therefore we have that \(A' = B' = A\) and from (2) and (4) we reconstruct \(C\) as follows:

n-adjacent vertices of \(\{x : x_n = \beta\}\) and \(\{x : x_n = \alpha'\}\) are both in \(C\) or not

(14)

Let all \(y \in \mathcal{A}^{n-1}\) such that \(y_i \in A, y_j \in \overline{A}\) be zeros of \(f\). Then from (13) we have that \(f(y) = 0\) when \(y_i \in A\) or \(y_j \in \overline{A}\) which implies that \(f\) is either the all-zero function or has two nonessential coordinates by Lemma 4. However, in the latter case a vertex \(y, y_i \in A, y_j \in \overline{A}\) is adjacent only to zeros of \(f\), which contradicts the definition of the cross. Therefore \(f\) is the all-zero function and we reconstruct \(C\) as follows:

n-adjacent vertices of \(\{x : x_n = \beta\}\) and \(\{x : x_n = \alpha\}\) are both in \(C\) or not

(15)

We now show that there are exactly \(q/2\) elements \(\beta \in A\) that satisfy (14).

Let \(z \in C\) be such that \(z_n = \alpha', z_i \in A, z_j \in \overline{A}\). Then by (3) the vertex \(z\) is not \(l\)-adjacent to vertices of \(\overline{C}\), \(l \in \mathcal{N} \setminus i, j, n\) and by (5) and (6) the vertex \(z\) is \(i\)- or \(j\)-adjacent to exactly \(q/2\) vertices of \(\overline{C}\). Therefore, \(z\) is \(n\)-adjacent with exactly \(q/2 = S_{12} - q/2\) vertices of \(\overline{C}\). In other words there are exactly \(q/2\) elements \(\beta \in A\) that satisfy (14).

Finally, from (8), (9), (14), (15) we see that \(x\) is in \(C\) iff

\[
(x_n, x_i, x_j, x_s) \in A \times A \times \overline{A} \times A \cup A \times A \times A \times \overline{A} \cup A \times A \times A \times A,
\]

i.e. the partition \((C, \overline{C})\) is obtained from Construction B.

The main result of this section is

Theorem 5. The only reduced equitable 2-partitions of \(H(n, q)\) with eigenvalue \(\lambda_2(n, q)\) are:

1. Reduced equitable 2-partitions of \(H(2, q)\) and \(H(3, q)\)
2. If \(q\) is even, the equitable 2-partition of \(H(4, q)\) from alphabet liftings of two induced 8-cycles in \(H(4, 2)\) (Construction B).
3. Equitable 2-partitions of \(H(n, q)\) obtained by the permutation switchings. (Construction A).

Proof. Consider \((\chi_C)^\alpha_n\), \(\alpha \in A\). Define the partition of \(A\) into sets \(A_1, \ldots, A_t\) that we call blocks as follows: \(\alpha, \beta\) are in one block iff \((\chi_C)^\alpha_n\) and \((\chi_C)^\beta_n\) have the same sets of essential coordinate positions. The proof follows from Lemmas (8) and (9) below.
Lemma 7. If there are at least two blocks, then for any $\alpha \in A$ the function $(\chi_C)_n^\alpha$ has exactly one essential coordinate.

Proof. Let $\alpha$ and $\beta$ be in different blocks. Then by Lemma (A) $(\chi_C)_n^\alpha - (\chi_C)_n^\beta$ is either the $(A, B, i)$-string or the $(A, B, i, j)$-cross.

The function $(\chi_C)_n^\alpha - (\chi_C)_n^\beta$ is not $(A, B, i)$-string because $\alpha$ and $\beta$ are in different blocks. Suppose the opposite. Obviously, $i$ is an essential coordinate for both $(\chi_C)_n^\alpha$ and $(\chi_C)_n^\beta$. Moreover, $i \in N \setminus \{i, n\}$ is not an essential coordinate for $(\chi_C)_n^\alpha - (\chi_C)_n^\beta$, so for any pair of $l$-adjacent tuples $y$ and $y'$ of $A^{n-1}$, we have that $(\chi_C)_n^\alpha(y) - (\chi_C)_n^\alpha(y') = (\chi_C)_n^\beta(y) - (\chi_C)_n^\beta(y')$. In other words, we have that $(\chi_C)_n^\alpha(y) - (\chi_C)_n^\alpha(y') = (\chi_C)_n^\alpha(y) - (\chi_C)_n^\alpha(y')$, so $i$ is an essential coordinate for both $(\chi_C)_n^\alpha$ and $(\chi_C)_n^\beta$ or not. We conclude that $(\chi_C)_n^\alpha$ and $(\chi_C)_n^\beta$ are in one block.

Let $(\chi_C)_n^\alpha - (\chi_C)_n^\beta$ be $(A, B, i, j)$-cross. If $(\chi_C)_n^\alpha$ or $(\chi_C)_n^\beta$ has an essential coordinate $s \in N \setminus \{i, j, n\}$, then by Theorem (A) we have Construction B and only one block in this case. So, we conclude that $(\chi_C)_n^\alpha$ and $(\chi_C)_n^\beta$ have exactly one essential coordinate in $\{i, j\}$.

Lemma 8. If there are at least two blocks then $(C, \overline{C})$ is obtained by Construction A.

Proof. By Lemma (A) the number of blocks is greater than the number of the essential coordinate positions of the partition by 1. Since the partition is reduced, the number of essential coordinates is $n$ and $A_1, \ldots, A_{n-1}$ are the blocks of the partition $(C, \overline{C})$. W.l.o.g. for any $\alpha \in A$, $i$ is the essential coordinate for $(\chi_C)_n^\alpha$.

Consider a vertex whose $n$th position is in $A_i$. Taking into account Proposition (A) and because each of $(\chi_C)_n^\alpha$ has exactly one essential coordinate by Lemma (14) $|C \cap K_i| = S_{21}/2$, $|\overline{C} \cap K_i| = S_{12}/2$ for the maximum clique $K_i$ consisting of the vertex and its $i$-neighbors. Let the vertex be from $C$ ($\overline{C}$ respectively). Then since $i$ is the only essential coordinate for $(\chi_C)_n^\alpha$, $\alpha \in A_i$, the vertex is not $j$-adjacent to any vertices of $\overline{C}$ ($C$ respectively) for $j \in N \setminus \{i, n\}$ and therefore is $n$-adjacent to exactly $S_{12}/2$ vertices of $\overline{C}$ ($S_{21}/2$ vertices of $C$). So we see that any maximum clique consisting of pairwise $n$-adjacent vertices of $H(n, q)$ contains exactly $S_{12}/2$ vertices of $\overline{C}$ and $S_{21}/2$ vertices of $C$.

Let $K_n$ and $K_n'$ be two maximum cliques consisting of pairwise $n$-adjacent vertices of $H(n, q)$. For $i \in N \setminus n$ denote by $L_i$ ($L_i'$) those vertices of $K_n$ and ($K_n'$ respectively) that have their $n$th coordinate in $A_i$. From the shown above we have that

$$|K_n \cap \overline{C}| = |K_n' \cap \overline{C}| = \sum_{i \in N \setminus n} |L_i \cap \overline{C}| = S_{12}/2. \quad (16)$$

We now prove that for any $i \in N \setminus n$ we have that

$$|L_i \cap \overline{C}| = |L_i' \cap \overline{C}| = S_{12}|L_i|/2q.$$
It is sufficient to prove the equality above when for some \( \beta \in A \) the vertices of \( K'_n \) is obtained from the vertices of \( K_n \) (and therefore \( L_j \) and \( L'_j \) for any \( j \in \mathcal{N} \setminus \{i, n\} \)) by changing their \( i \)-th coordinate to \( \beta \). We see that \( i \)-adjacent pairs of vertices between \( L_j \) and \( L'_j \) is a perfect matching. This, combined with the fact that \( i \) is not essential for the restriction of \( (C, \overline{G}) \) to \( A^{n-1} \times A_j \) gives \( |L_j \cap \overline{C}| = |L'_j \cap \overline{C}| \). Then from (15) we have that \( |L_i \cap \overline{C}| = |L'_i \cap \overline{C}| \).

For \( j \in \mathcal{N} \setminus n \) consider the restriction of the partition \((C, \overline{G})\) to the subgraph induced by \( A^{n-1} \times A_j \). Each maximum clique of the subgraph consisting of pairwise \( j \)-adjacent vertices contains \( S_{12}/2 \) vertices of \( \overline{G} \) and any maximal clique of the subgraph consisting of pairwise \( n \)-adjacent vertices contains has exactly \( l_j \) vertices of \( \overline{G} \). Since the coordinates from \( \mathcal{N} \setminus \{j, n\} \) are nonessential for the restriction, double counting of \( \overline{G} \) in \( A^{n-1} \times A_j \) gives that \( l_j/|A_j| = S_{12}/2q \) is a constant regardless of \( j \). We have shown that the partition \((C, \overline{G})\) is equally distributed by all maximum cliques \( K_j, j \in \mathcal{N} \setminus n \), whose vertices have \( n \)-th coordinate in \( A_j \) and by the sets \( x' \times A_j \), for all \( x' \in A^{n-1}, j \in \mathcal{N} \). We conclude that \((C, \overline{G})\) is obtained by Construction \( A \) by the transposition of coordinate positions 1 and \( n \).

\[ \square \]

**Lemma 9.** If there is one block then \((C, \overline{G})\) has not more than 3 essential coordinates or \((C, \overline{G})\) is obtained by Construction \( B \).

**Proof.** Let \( \alpha, \beta \) be distinct elements of \( A \). Consider \((\chi_C)^\alpha_n - (\chi_C)^\beta_n\). By Lemma 4 and Theorem 4 we have the following cases:

- \((\chi_C)^\alpha_n - (\chi_C)^\beta_n\) is the \((A, B, i, j)\)-cross and there is \( s \in \mathcal{N} \setminus \{i, j, n\} \) that is essential for \((\chi_C)^\alpha_n\) and \((\chi_C)^\beta_n\). Then \((C, \overline{G})\) is obtained by Construction \( B \).

- \((\chi_C)^\alpha_n - (\chi_C)^\beta_n\) is the \((A, B, i, j)\)-cross, the essential coordinates of \((\chi_C)^\alpha_n\) and \((\chi_C)^\beta_n\) are in \( \{i, j\} \). Since \( \alpha \) and \( \beta \) are from the same block, \( i \) and \( j \) are the essential coordinates for both \((\chi_C)^\alpha_n\) and \((\chi_C)^\beta_n\). The partition \((C, \overline{G})\) has exactly three essential coordinates in this case.

- \((\chi_C)^\alpha_n - (\chi_C)^\beta_n\) is the \((A, B, i)\)-string.

Let \((\chi_C)^\alpha_n - (\chi_C)^\beta_n\) be the \((A, B, i)\)-string. From the definition of the string we have that \( x \in C \) if \( x_n = \alpha, x_i \in A \) or \( x_n = \beta, x_i \in B \) and \( x \in \overline{C} \) if \( x_n = \beta, x_i \in A \) or \( x_n = \alpha, x_i \in B \).

Suppose that, \( s \neq i, n \) is an essential coordinate of \((\chi_C)^\alpha_n\) and \((\chi_C)^\beta_n\) and there are \( s \)-adjacent vertices for \( x_{000} \in C \) and \( x_{001} \in \overline{C} \), \( x_{000} = x_{001} = \alpha \), \( x_{000} = x_{001} = \beta \). Fix \( \alpha' \in A \) and denote by \( x^a \) the vertex which is obtained from \( x_{000} \) by changing its \( n \)-th position to \( \beta \) iff \( a_1 = 1 \) and \( i \)-th position to \( \alpha' \) iff \( a_2 = 1 \) and \( s \)-th position to \( x_{001} = a_3 = 1 \). By the properties above we know that \( x^a \in C \) iff \( a \in \{000, 010, 011, 100\} \). Consider \( f = (\chi_C)^\beta_n - (\chi_C)^\alpha_n \) on the tuples \( \{x^{a_1a_2a_3} : a_1, a_3 \in \{0, 1\} \} \), here \( x^{a_1a_2a_3} \) is obtained from \( x^{a_1a_2a_3} \) by deleting its \( i \)-th position. We have that \( f(x^{000}) = 0 \), \( f(x^{011}) = -1 \), \( f(x^{10}) = 1 \), which implies that \( f \) is a cross and we are in the second case.

\[ \square \]
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