Alternating-time temporal logic with resource bounds

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Abstract

Many problems in AI and multi-agent systems research are most naturally formulated in terms of the abilities of a coalition of agents. There exist several excellent logical tools for reasoning about coalitional ability. However, coalitional ability can be affected by the availability of resources, and there is no straightforward way of reasoning about resource requirements in logics such as Coalition Logic (CL) and Alternating-time Temporal Logic (ATL). In this article, we describe a logic for reasoning about coalitional ability under resource constraints. We extend ATL with costs of actions and hence of strategies. We give a complete and sound axiomatization of the resulting logic, Resource-Bounded ATL (RB-ATL) and a model-checking algorithm for it.

Keywords: Strategic ability, resources, axiomatisation of RB-ATL.

1 Introduction

In many situations, a group of agents can cooperate to achieve an outcome which cannot be achieved by any agent in the group acting individually. For example, in the prisoners dilemma, a single prisoner cannot ensure the optimal outcome, while a coalition of two prisoners can. Similarly, it may be possible for a set of cooperating agents to solve a difficult computational problem by distributing it, while a single agent may not have sufficient memory or processor power to solve it. In the latter case, there is an interaction between the amount of resources available to the agents (or the amount of resources which they are willing to contribute), and their ability to jointly achieve a goal.

In this article we describe a logic, Resource-Bounded Alternating-time Temporal Logic (RB-ATL), for reasoning about coalitional ability under resource constraints. RB-ATL allows us to express and verify properties such as:

1. a coalition of agents $A$ has a strategy to achieve a property $\varphi$ provided they have resources $b$, but they cannot enforce $\varphi$ under a tighter resource bound $b_1$;
2. $A$ has a strategy to maintain the property $\varphi$, provided they have resources $b'$; and
3. $A$ has a strategy to maintain $\varphi$ until $\psi$ becomes true, provided $A$ has resources $b'$.

In Section 2.4, we illustrate the expressive power of RB-ATL on a simple example of a sensor network, where the agents (sensor nodes) require two resources: energy and memory.

In previous work, we studied a version of Coalition Logic (CL) with resource bounds, RBCL [5]. RBCL can express properties of the form (1) above, but not of the form (2) and (3). Other work on

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temporal logics and logics of coalitional ability with resource constraints includes \cite{2, 7, 8, 10, 11}. However, this work concentrates on model-checking complexity, rather than axiomatization, which is the focus of this article.

This article is a revised and extended version of \cite{4}. In \cite{4}, we gave a sound and complete axiomatization and a model-checking algorithm for a version of RB-ATL without infinite resource bounds. The main differences from \cite{4} are the addition of an infinite resource bound (to make the logic a conservative extension of ATL), and the addition of complete proofs and an illustrative example.\footnote{A preliminary version of RB-ATL with infinite bounds was introduced in \cite{14}.}

The remainder of this article is organized as follows. In Section 2, we present the syntax and semantics of RB-ATL and show how RB-ATL can be used to express properties of a simple sensor network. In Section 3 we provide a sound and complete axiomatization of RB-ATL. In Section 4, we give a model-checking algorithm for RB-ATL. Finally, we survey related work in Section 5 and conclude in Section 6.

## 2 Syntax and semantics of RB-ATL

Consider a system of agents that can perform actions to change the state (we assume concurrent execution of actions by all agents). We denote the set of agents by $N$. To reason about resources, we assume that actions have costs. Let $R$ be a set of resources (such as money, energy, or anything else which may be required by an agent for performing an action). We assume that a cost of an action, for each of the resources, is a non-negative integer. The set of resource bounds $\mathbb{B}$ over $R$ is defined as $\mathbb{B} = (\mathbb{N} \cup \{\infty\})^r$, where $r = |R|$. We denote by $\bar{0}$ the smallest resource bound $(0, \ldots, 0)$ and $\bar{\infty}$ the greatest resource bound $(\infty, \ldots, \infty)$.

### 2.1 Syntax of RB-ATL

The syntax of RB-ATL is defined as follows, where $A$ is a non-empty subset of $N$ and $b \in \mathbb{B}$.

$$\varphi ::= p | \neg \varphi | \varphi \lor \psi | \langle \langle A^b \rangle \rangle \bigcirc \varphi | \langle \langle A^b \rangle \rangle \Box \varphi | \langle \langle A^b \rangle \rangle \varphi U \psi$$

Here, $\langle \langle A^b \rangle \rangle \bigcirc \varphi$ means that a coalition $A$ can ensure that the next state satisfies $\varphi$ under resource bound $b$. $\langle \langle A^b \rangle \rangle \Box \varphi$ means that $A$ has a strategy to make sure that $\varphi$ is always true, and the cost of this strategy is at most $b$. Similarly, $\langle \langle A^b \rangle \rangle \varphi U \psi$ means that $A$ has a strategy to enforce $\psi$ while maintaining the truth of $\varphi$, and the cost of this strategy is at most $b$. Notice the meaning of these operators when $b = \bar{\infty}$ is the same as their counterparts in ATL; in other words, the ATL operator $\langle \langle A \rangle \rangle$ corresponds to $\langle \langle A^\infty \rangle \rangle$ for $A \neq \emptyset$ and $\langle \langle \emptyset \rangle \rangle$ to the dual of $\langle \langle N^\infty \rangle \rangle$.

### 2.2 Semantics of RB-ATL

To interpret this language, we extend the definition of concurrent game structures \cite{6} with resource requirements for executing actions. For consistency with \cite{6}, in what follows we refer to agents as ‘players’ and actions as ‘moves’.

**Definition 1**

A Resource-bounded Concurrent Game Structure (RB-CGS) is a tuple $S = (n, r, Q, \Pi, \pi, d, c, \delta)$ where:

- $n \geq 1$ is the number of players (agents), we denote the set of players $\{1, \ldots, n\}$ by $N$;
We assume that each agent in each state has an available action with \( m_A \) option of doing nothing.

Given a RB-CGS \( Q \), \( Q \) set of finite non-empty sequences of states by

\[ Q = \{q_1, q_2, \ldots\} \]

We assume that each agent in each state has an available action with \( 0 \) cost (intuitively, it has the option of doing nothing).

Given a RB-CGS \( S \), we denote by \( Q^* \) the set of finite sequences of states or finite computations and by \( Q^\omega \) the set of infinite sequences of states or infinite computations. For a finite or infinite computation \( \lambda = q_1 q_2 \ldots \in Q^* \cup Q^\omega \), we use the notation \( \lambda[i] = q_i \) and \( \lambda[i, j] = q_i \ldots q_j \). We denote the set of finite non-empty sequences of states by \( Q^+ \).

**Definition 2**

Given a RB-CGS \( S \) and a state \( q \in Q \), a move (or a joint action) for a coalition \( A \subseteq N \) is a tuple \( \sigma_A = (\sigma_a)_{a \in A} \) such that \( 1 \leq \sigma_a \leq d(q, a) \).

By \( D_A(q) \) we denote the set of all moves for \( A \) at state \( q \). Given a move \( m \in D(q) \), we denote by \( m_A \) the actions executed by \( A \), \( m_A = (m_a)_{a \in A} \). We define the set of all possible outcomes of a move \( \sigma_A \in D_A(q) \) at state \( q \) as follows:

\[ \text{out}(q, \sigma_A) = \{q' \in Q \mid \exists m \in D(q) : m_A = \sigma_A \land q' = \delta(q, m)\} \]

For convenience, we define the projection of \( \infty \) components of a resource bound \( b \) on another bound \( d \) as \( d \xrightarrow{\infty} b \) where for all \( i \in [1, \ldots, r] \):

\[ (d \xrightarrow{\infty} b)_i = \begin{cases} d_i & \text{if } b_i \neq \infty \\ \infty & \text{if } b_i = \infty \end{cases} \]

For example, \((2, \infty, \infty, 6) \xrightarrow{\infty} (1, \infty, 3, \infty) = (2, \infty, \infty, \infty)\) and \(0 \xrightarrow{\infty} (1, \infty, 3, \infty) = (0, \infty, 0, \infty)\). To compare costs and resource bounds, we use the usual pointwise vector comparison, i.e. \((b_1, \ldots, b_r) \leq (d_1, \ldots, d_r)\) iff \( b_i \leq d_i \) for \( i \in [1, \ldots, r] \) where \( n \leq \infty \) for all \( n \in \mathbb{N} \). We also use pointwise vector addition: \((b_1, \ldots, b_r) + (d_1, \ldots, d_r) = (b_1 + d_1, \ldots, b_r + d_r)\) where \( n + \infty = \infty \) for all \( n \in \mathbb{N} \cup \{\infty\} \). Conversely, we also split a resource bound \( b \) into pairs of resource bounds \((d, d')\) such that:

(i) \( d + d' = b \),
(ii) \( d_i = d_i' = \infty \) for all \( i \in [1, \ldots, r] \) such that \( b_i = \infty \), and
(iii) \( d \neq 0 \xrightarrow{\infty} b \).
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The set of all such pairs \((d, d')\) is denoted by \(\text{split}(b)\). Obviously, \(\text{split}(b)\) is finite.

The cost of a move \(\sigma A \in D_A(q)\) is defined as \(\text{cost}(q, \sigma A) = \Sigma_{a \in A} (q, a, \sigma A)\). (Note that we use \(c\) for the cost of single actions and \(\text{cost}\) for the cost of joint actions.)

**Definition 3**

Given a RB-CGS \(S\), a strategy for a subset of players \(A \subseteq N\) is a mapping \(F_A\) which associates each sequence \(\lambda q \in Q^+\) to a move in \(D_A(q)\).

A computation \(\lambda \in Q^\omega\) is consistent with \(F_A\) iff for all \(i \geq 1\), \(\lambda[i + 1] \in \text{out}(\lambda[i], F_A(\lambda[1..i]))\). We denote by \(\text{out}(q, F_A)\) the set of all such sequences \(\lambda\) starting from \(q\), i.e. \(\lambda[1] = q\). Given a non-empty finite prefix \(\lambda\) of a computation which is consistent with a strategy \(F_A\), we define the cost of \(F_A\) with respect to \(\lambda\) as \(\text{cost}(\lambda, F_A) = \sum_{i=1}^{\lambda - 1} \text{cost}(\lambda[i], F_A(\lambda[1..i]))\).

**Definition 4**

Given a bound \(b\), a computation \(\lambda \in \text{out}(q, F_A)\) is \(b\)-consistent with \(F_A\) iff, for every finite prefix \(\lambda'\) of \(\lambda\), \(\lambda' \preceq_b \lambda\). We denote by \(\text{out}(q, F_A, b)\) the set of all \(b\)-consistent computations. A strategy \(F_A\) is a \(b\)-strategy iff \(\text{out}(q, F_A) = \text{out}(q, F_A, b)\) for any \(q \in Q\).

In other words, all executions of a \(b\)-strategy cost at most \(b\) resources. Note that this means that each computation of such a strategy starts with a finite prefix where some non-\(0 \preceq b\) cost actions are executed, and continues with an infinite sequence of \((0 \preceq b)\)-cost actions.

2.3 Truth definition for RB-ATL

Given a RB-CGS \(S = (n, r, Q, \Pi, \pi, d, c, \delta)\), the truth definition for RB-ATL is given inductively as follows:

- \(S, q \models \psi\) if and only if \(q \models \pi\);
- \(S, q \models \psi\) if and only if \(q \not\models \psi\);
- \(S, q \models \psi \lor \psi\) if and only if \(S, q \models \psi\) or \(S, q \models \psi\);
- \(S, q \models (\langle A \rangle) \psi\) if there exists a \(b\)-strategy \(F_A\) such that for all \(\lambda \in \text{out}(q, F_A)\), \(S, \lambda[2] \models \psi\) if there is a move \(\sigma A \in D_A(q)\) such that for all \(\lambda' \in \text{out}(\sigma A)\), \(S, q' \models \psi\);
- \(S, q \models (\langle A \rangle) \psi\) if there exists a \(b\)-strategy \(F_A\) for any \(\lambda \in \text{out}(q, F_A)\), \(S, \lambda[i] \models \psi\) for all \(i \geq 1\);
- \(S, q \models (\langle A \rangle) \psi\) if there exists a \(b\)-strategy \(F_A\) such that for all \(\lambda \in \text{out}(q, F_A)\), there exists \(i \geq 1\) such that \(S, \lambda[i] \models \psi\) and \(S, \lambda[j] \models \psi\) for all \(j \in \{1, \ldots, i - 1\}\).

Notice that the truth definition of \((\langle A \rangle)\) is the same as that of \((\langle A \rangle)\) in ATL.

2.4 Example

To conclude this section, we describe a concrete scenario to illustrate the notions introduced above, and give some examples of the expressive power of RB-ATL.

Consider a sensor network consisting of two agents (sensor nodes), 1 and 2. The agents monitor for movement. If they detect movement, they can inform their neighbour. If an agent receives a communication from its neighbour, it can save it. If an agent has more than one record of movement, the agent can report this to the base station. We assume that 2 is closer to the base station than 1. We consider two resources, energy and memory. Sending a message requires energy (depending on the distance to the receiver) and saving a communication requires memory. Sending from 1 to 2 (\(\text{send}12\)) and from 2 to 1 (\(\text{send}21\)) both require 2 units of energy and 0 memory. Saving a record requires 0...
units of energy and 1 unit of memory. Sending from 1 to the base station \((send1b)\) requires 3 units of energy, and sending from 2 to the base \((send2b)\) requires 1 unit of energy. The option of doing nothing \((idle)\) is always available and costs nothing. In the initial state \(q_0\), each agent has a record of having itself seen movement. The system is shown in Figure 1, where transitions between states are annotated with tuples of actions (the first element is an action by agent 1, and the second is an action by agent 2). We omit self-loops in each state by the joint action \((idle, idle)\) for readability.

In this scenario, both agents together can enforce the outcome \(q_6\), where we assume that a proposition \(p\) (which means that the base station has been informed) holds. Moreover, they can achieve this by spending 3 units of energy and 1 unit of memory by choosing the following actions: \((send12, idle)\) in \(q_0\), \((idle, save)\) in \(q_1\) and \((idle, send2b)\) in \(q_4\). This can be expressed in RB-ATL as \(\langle\langle\{1, 2\}\rangle\rangle^{\leq 3, 1} U p\). It is also the case that the agents cannot achieve this without using some memory, even if they use unlimited energy: \(\neg\langle\langle\{1, 2\}\rangle\rangle^{\leq (\infty, 0)} U p\). Clearly, neither of the agents can single-handedly enforce \(q_6\); however, once the system is in \(q_6\), either agent can trivially maintain \(p\) forever, since the only choice of action available to each agent there is \(idle\). This can be expressed as \(\langle\langle\{1, 2\}\rangle\rangle^{\leq (3, 1)} U \langle\langle\{1\}\rangle\rangle^0 \square p\).

### 3 Axiomatization

In this section we present the axiomatic system for RB-ATL. To make the formulas below more readable, we define the following abbreviations:

\[
\begin{align*}
\langle\langle A^b \rangle\rangle \bigcirc \varphi & = \bigvee_{(d, d') \in \text{split}(b)} \langle\langle A^d \rangle\rangle \bigcirc \langle\langle A^{d'} \rangle\rangle \square \varphi \\
\langle\langle A^b \rangle\rangle \bigcirc \varphi \bigcup \psi & = \bigvee_{(d, d') \in \text{split}(b)} \langle\langle A^d \rangle\rangle \bigcirc \langle\langle A^{d'} \rangle\rangle \varphi \bigcup \psi
\end{align*}
\]

The axiomatic system consists of the following axiom schemas and rules of inference, where \(A\), \(A_1\) and \(A_2\) are non-empty subsets of \(N\), and \(b, d \in \mathbb{N}\).

#### Axioms:

- **(PL)** Tautologies of Propositional Logic
- **(1)** \(\neg\langle\langle A^b \rangle\rangle \bigcirc \bot\)
- **(T)** \(\langle\langle A^b \rangle\rangle \bigcirc \top\)
- **(B)** \(\langle\langle A^b \rangle\rangle \bigcirc \varphi \rightarrow \langle\langle A^d \rangle\rangle \bigcirc \varphi\)

where \(b \leq d\)
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\begin{align*}
\text{(S)} & \quad \langle\langle A_1^b \rangle\rangle \circ \varphi \land \langle\langle A_2^d \rangle\rangle \circ \psi \rightarrow \langle\langle A_1 \cup A_2^{b+d} \rangle\rangle \circ (\varphi \land \psi) \quad \text{where } A_1 \cap A_2 = \emptyset \\
\text{(SN)} & \quad \langle\langle A^b \rangle\rangle \circ \varphi \land \neg \langle\langle N^\infty \rangle\rangle \circ \neg \psi \rightarrow \langle\langle A^b \rangle\rangle \circ (\varphi \land \psi) \\
\text{(SN+)} & \quad \langle\langle N^b \rangle\rangle \circ \varphi \land \neg \langle\langle N^b \rangle\rangle \circ \neg \psi \rightarrow \langle\langle N^b \rangle\rangle \circ (\varphi \land \psi) \\
\text{(SN−)} & \quad \neg \langle\langle N^b \rangle\rangle \circ \neg \varphi \land \neg \langle\langle N^b \rangle\rangle \circ \neg \psi \rightarrow \neg \langle\langle N^b \rangle\rangle \circ \neg (\varphi \land \psi) \\
\text{(FPc)} & \quad \langle\langle A^b \rangle\rangle \circ \varphi \land \neg \langle\langle A^b \rangle\rangle \circ \sigma \circ \varphi \lor \langle\langle 0^\infty \rangle\rangle \circ (\langle\langle A^b \rangle\rangle \circ \varphi) \\
\text{(FPu)} & \quad \langle\langle A^b \rangle\rangle \psi U \psi \leftrightarrow \psi \land (\langle\langle A^b \rangle\rangle \circ \psi U \psi \lor \langle\langle 0^\infty \rangle\rangle \circ (\langle\langle A^b \rangle\rangle \psi U \psi))
\end{align*}

Inference rules:

\begin{align*}
\text{(MP)} & \quad \varphi, \varphi \rightarrow \psi \\
\text{\langle\langle A^b \rangle\rangle \circ-Monotonicity} & \quad \varphi \rightarrow \psi \\
\text{\langle\langle N^b \rangle\rangle \circ-Necessitation} & \quad \varphi \rightarrow \langle\langle N^b \rangle\rangle \circ \neg \varphi \\
\text{\langle\langle A^b \rangle\rangle \circ-Induction} & \quad \sigma \rightarrow (\varphi \land (\langle\langle A^b \rangle\rangle \circ \sigma \rightarrow (\langle\langle A^b \rangle\rangle \circ \circ \varphi \lor \langle\langle 0^\infty \rangle\rangle \circ \theta))) \\
\text{\langle\langle A^b \rangle\rangle U-Induction} & \quad (\varphi \land (\langle\langle A^b \rangle\rangle \circ \varphi U \psi \lor \langle\langle 0^\infty \rangle\rangle \circ (\langle\langle A^b \rangle\rangle \psi U \psi))) \rightarrow \theta \\
\end{align*}

Before proving soundness and completeness, i.e., every formula derived by the above system is valid and every valid formula can be derived by the above system, we give an intuitive explanation of the axioms and compare them with the axiomatic system for ATL given in [13].

First of all, observe that with the resource bounds removed, the axioms (., (.), (S), and the inference rules (\langle\langle A^b \rangle\rangle \circ-Monotonicity) and (\langle\langle N^b \rangle\rangle \circ-Necessitation) are identical to their ATL counterparts. Unlike ATL, we need several versions of (S) since we do not have the (\langle\langle \theta^b \rangle\rangle) modality and, as a result, (SN), (SN+) and (SN−) are not derivable from (S). The axiom (B) says that if A can enforce \varphi under a resource bound b, then it can also enforce \varphi if it has more than b resources. The axiom (FPc) is similar to its ATL counterpart. However, unlike in ATL, there are two ways to ‘unwind’ \langle\langle A^b \rangle\rangle \circ \psi in RB-ATL: one way is to make a move which costs a non-trivial amount of resources d, and then maintain \varphi with \b - d resources; the second way is to make a trivial (\emptyset \infty b)-cost move, and then maintain \varphi with b resources. Similarly for (FPu). Finally, the rules (\langle\langle A^b \rangle\rangle \circ-Induction) and (\langle\langle A^b \rangle\rangle U-Induction) correspond to the ATL axioms (GFPc) and (LFPu); the first one says that \circ corresponds to the greatest fixed point and the second that U corresponds to the least fixed point.

This will be made more precise after we give fixed point characterizations of the temporal operators.

3.1 Fixed point characterizations of temporal operators

Consider an operation \([\langle\langle A^b \rangle\rangle \circ]\), which, given a set of states X, returns the set of states from where A can enforce an outcome to be in X under resource bound b (this is the same as Pre(A,X,b) defined in Section 4, which is in turn similar to Pre from [6]):

**Definition 5**

\([\langle\langle A^b \rangle\rangle \circ]: \varphi(Q) \rightarrow \varphi(Q) \text{ is defined as follows: given a set } X \subseteq Q, [\langle\langle A^b \rangle\rangle \circ](X) \text{ is the set}

\[\{q \mid \exists \sigma \in D_A(q): \text{cost}(q,\sigma) \leq b \land \text{out}(q,\sigma) \subseteq X\}\]
Let us define \( \| \varphi \| = \{ q \in Q | S, q \models \varphi \} \). It is straightforward that:

\[
\| (\langle A^b \rangle) \varphi \| = \| (\langle A^b \rangle) \Box \| \varphi \|
\]

Recall that if \( f \) is a monotone operator \( 2^Q \to 2^Q \) (that is, \( X \subseteq Y \) implies \( f(X) \subseteq f(Y) \)), then \( X \) is a fixed point of \( f \) if \( f(X) = X \). By the Knaster–Tarski theorem, \( f \) has the least and the greatest fixed point. The least fixed point of \( f \) is denoted by \( \mu X f(X) \) and the greatest fixed point by \( \nu X f(X) \). We are going to show that the meanings of \( \Box \) and \( \mathcal{U} \) correspond to the greatest and the least fixed points of certain operations on sets of states.

**Lemma 1**

For all \( q \in Q \), the following fixed point characterizations hold:

1. \( q \in \| (\langle A^b \rangle) \Box \| \varphi \| \) iff \( q \in v X \). \( \| \varphi \| \cap \| (\langle A^b \rangle) \Box \| \varphi \| \cup \| (\langle A^{\infty} \rangle) \Box \| \varphi \| \) iff there is a \( b \)-strategy \( F_A \) for \( A \) such that for all \( \lambda \in out(q, F_A) \), \( \lambda[i] \in \| \varphi \| \) for all \( i \geq 1 \).

2. \( q \in \| (\langle A^b \rangle) \mathcal{U} \| \varphi \| \) iff \( q \in \mu X \). \( \| \varphi \| \cup \| (\langle A^b \rangle) \mathcal{U} \| \varphi \| \cup \| (\langle A^{\infty} \rangle) \Box \| \varphi \| \) iff there is a \( b \)-strategy \( F_A \) for \( A \) such that for all \( \lambda \in out(q, F_A) \), there exists \( i \geq 1 \) such that \( \lambda[i] \in \| \varphi \| \) and \( \lambda[j] \in \| \varphi \| \) for all \( j < i \).

**Proof.**

We will only provide the proof for the first case as the second can be done in a similar way. For convenience, let us denote \( f(X) = \| \varphi \| \cap \| (\langle A^b \rangle) \Box \| \varphi \| \cup \| (\langle A^{\infty} \rangle) \Box \| \varphi \| \) for \( \langle A^b \rangle \mathcal{U} \lambda \varphi \| \). We will only provide the proof for the first case as the second can be done in a similar way.

To show that \( \| (\langle A^b \rangle) \Box \| \varphi \| \) is a post-fixed point of \( f(X) \), i.e. \( f(Y) \supseteq Y \). Let \( q \in Y \), by the semantics definition, we have that there is a \( b \)-strategy \( F_A \) such that for any \( \lambda \in out(q, F_A) \), \( \lambda[i] \in \| \varphi \| \) for all \( i \geq 1 \). Then, \( q = \lambda[1] \in \| \varphi \| \). Assume that \( b^* = \text{cost}(q, F_A(q)) \), let \( b'' \) be a resource bound such that \( (b'' \in \infty, b'') \in \text{split}(b) \). For each \( q' \in out(q, F_A(q)) \), we define a \( b'' \)-strategy \( F_{q'}^* A \) as the remainder of \( F_A \) from \( q' \), i.e. \( F_{q'}^* A(q') = F_A(q (k) k) \) for all \( k \in Q^n \). Then, for all \( q' \in out(q, F_{q'}^* A) \), we have that \( q'' \leq \infty \). It is straightforward that any computation in \( out(q', F_{q'}^* A) \) costs at most \( b'' \). Then, \( q'' \in \| (\langle A^b \rangle) \Box \| \varphi \| \).

Thus, \( q \in \| (\langle A^b \rangle) \Box \| \varphi \| \). If \( b'' < \infty \), we have that \( q \in \| (\langle A^b \rangle) \Box \| \varphi \| \). Otherwise \( q \in \| (\langle A^{\infty} \rangle) \Box \| \varphi \| \). This means that \( q \in f(Y) \).

To show that \( \| (\langle A^b \rangle) \Box \| \varphi \| \) is in fact the greatest fixed point of \( f(X) \), we show that, for every post-fixed point \( Z, Z \subseteq Y \).

We have \( f(X) = \| \varphi \| \cap \| (\langle A^b \rangle) \Box \| \varphi \| \cup \| (\langle A^{\infty} \rangle) \Box \| \varphi \| \). Assume \( q \in Z \), we have:

\[
q \in Z \Rightarrow q \in \| \varphi \| \text{ and either } q \in \| (\langle A^b \rangle) \Box \| \varphi \| \text{ or } q \in \| (\langle A^{\infty} \rangle) \Box \| \varphi \| \text{ (Z).}
\]

We define a \( b \)-strategy \( F_A \) which will maintain \( \varphi \) by induction on the length of inputs for \( F_A \). Let \( A^i \) denote the set of inputs of length \( i \geq 1 \) for \( F_A \). Initially, \( A^1 = \{ q \} \). We will define \( F_A \) for input of length \( i \) and \( A^{i+1} \) inductively on \( i \geq 1 \) such that, for all \( \lambda \in A^{i+1} \), either \( \text{cost}(\lambda, F_A) \leq d \) and \( \lambda[i+1] \in \| (\langle A^b \rangle) \Box \| \varphi \| \) for some \( (d, d') \in \text{split}(b) \) or \( \text{cost}(\lambda, F_A) \leq \infty \) and \( \lambda[i+1] \in Z \).

- Case \( i = 1 \), recall that \( q \in \| \varphi \| \) and either \( q \in \| (\langle A^b \rangle) \Box \| \varphi \| \) or \( q \in \| (\langle A^{\infty} \rangle) \Box \| \varphi \| \).
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- If \( q \in [(\langle A^d \rangle) \circ \Box \varphi] \), then there exists \((d, d') \in \text{split}(b)\) such that \( q \in [(\langle A^d \rangle) \circ (\langle A^{d'} \rangle) \Box \varphi] \).

  then, we have:

  \[
  q \in [(\langle A^d \rangle) \circ (\langle A^{d'} \rangle) \Box \varphi] \Rightarrow q \in [(\langle A^d \rangle) \circ ([\langle A^{d'} \rangle] \Box \varphi)]
  \]

  \[
  \Rightarrow \exists \sigma_a \in D_A(q) : \text{cost}(q, \sigma_a) \leq d \land \\
  \hspace{1cm} \text{out}(q, \sigma_a) \subseteq [(\langle A^d \rangle) \Box \varphi]
  \]

  \[
  \Rightarrow \forall q' \in \text{out}(q, \sigma_a), \exists d'-\text{strategy } F_{A, q'}.
  \]

  \[
  \forall \lambda \in \text{out}(q', F_{A, q'}). \forall j \geq 1 : \lambda[j] \in \varphi
  \]

  Then, we define \( F_A(q) = \sigma_a \) and \( \Lambda^2 = \{qq' | q' \in \text{out}(q, \sigma_a)\} \). Obviously, we have:

  \[
  \forall qq' \in \Lambda^2 : \text{cost}(qq', F_A) = \text{cost}(q, \sigma_a) \leq d \land q \in [(\langle A^{d'} \rangle) \Box \varphi]
  \]

- If \( q \in [(\langle A^{\bar{d}-b} \rangle) \circ \Box \varphi] \), we have:

  \[
  q \in [(\langle A^{\bar{d}-b} \rangle) \circ \Box \varphi] \Rightarrow \exists \sigma_a \in D_A(q) : \text{cost}(q, \sigma_a) \leq \bar{d} \land \\
  \hspace{1cm} \text{out}(q, \sigma_a) \subseteq Z
  \]

  Let \( F_A(q) = \sigma_a \) and \( \Lambda^2 = \{qq' | q' \in \text{out}(q, \sigma_a)\} \). Then, we have:

  \[
  \forall qq' \in \Lambda^2 \Rightarrow \text{cost}(qq', F_A) = \text{cost}(q, \sigma_a) \leq \bar{d} \land q \in \text{out}(q, \sigma_a)
  \]

  \[
  \Rightarrow q' \in Z \text{ as } \text{out}(q, \sigma_a) \subseteq Z
  \]

• Case \( i > 1 \), let us assume that \( F_A \) for inputs of length \( i - 1 \) and \( \Lambda^i \) have been defined. By the induction hypothesis, we have, for all \( \lambda \in \Lambda^i \), either \( \text{cost}(\lambda, F_A) \leq d \) and \( \lambda[i] \in [(\langle A^d \rangle) \Box \varphi] \) for some \((d, d') \in \text{split}(b)\) or \( \text{cost}(\lambda, F_A) \leq \bar{d} \land \lambda[i] \in Z \). Let us define \( F_A \) for each input \( \lambda \in \Lambda^i \) and \( \Lambda^{i+1} \).

  - If \( \text{cost}(\lambda, F_A) \leq d \) and \( \lambda[i] \in [(\langle A^d \rangle) \Box \varphi] \), we have:

    \[
    \lambda[i] \in [(\langle A^d \rangle) \Box \varphi] \Rightarrow \exists d'-\text{strategy } F_{A, \lambda}, \\
    \forall \lambda' \in \text{out}(\lambda[i], F_{A, \lambda}), \\
    \forall j \geq 1 : \lambda'[j] \in \varphi
    \]

    Then, we define \( F_A(\lambda) = F_{A, \lambda}(\lambda[i]) \) and \( \Lambda^{i+1} = \{q | \lambda \in \Lambda^i \land q \in \text{out}(\lambda[i], F_{A, \lambda}(\lambda[i]))\} \).

    Let \( \text{cost}(\lambda[i], F_{A, \lambda}(\lambda[i])) = d'' \), we have, for all \( \lambda q' \in \Lambda^{i+1} \) (where \( \lambda \in \Lambda^i \) and \( q' \in \text{out}(\lambda[i], F_{A, \lambda}(\lambda[i]))\):

    \[
    \text{cost}(\lambda q', F_A) = \text{cost}(\lambda, F_A) + \text{cost}(\lambda[i], F_{A, \lambda}(\lambda[i])) \leq d + \text{cost}(\lambda[i], F_{A, \lambda}(\lambda[i])) = d + d''
    \]

    Furthermore, by considering the \((d' - d'')\)-strategy \( F_{A, \lambda q'} \) where \( F_{A, \lambda q'}(\lambda') = F_{A, \lambda}(q' \lambda') \) for all \( \lambda' \in \mathcal{Q}^+ \), we have:

    \[
    \forall \lambda' \in \text{out}(q', F_{A, \lambda q'}). \forall j \geq 1 : \lambda'[j] \in \varphi \Rightarrow q' \in [(\langle A^{d'-d''} \rangle) \Box \varphi]
    \]

    Finally, it is straightforward that \(((d + d''), (d' - d'')) \in \text{split}(b)\).
If \( \lambda[i] \in Z \) and \( \lambda'[i] \in Z \), we have

\[
\lambda[i] \in Z \Rightarrow \lambda'[i] \in f(Z)
\]

\[
\Rightarrow \lambda'[i] \in \| \varphi \| \quad \text{and either } \lambda[i] \in \| (\langle A^b \rangle \bigcirc □ \varphi) \| \quad \text{or}
\]

\[
\lambda[i] \in \| (\langle A^0 \rangle \bigcirc A) \bigcirc (\langle Z \rangle) \|
\]

* If \( \lambda[i] \in \| (\langle A^b \rangle \bigcirc □ \varphi) \| \), then there exists \( (d,d') \in \text{split}(b) \) such that \( \lambda[i] \in \| (\langle A^d \rangle \bigcirc □ \varphi) \| \); then, we have:

\[
\lambda[i] \in \| (\langle A^d \rangle \bigcirc (\langle A^d \rangle \bigcirc □ \varphi) \|
\]

\[
\Rightarrow \exists \sigma_A \in D_A(\lambda[i]): \text{cost}(\lambda[i], \sigma_A) \leq d \quad \text{and}
\]

\[
\text{out}(\lambda[i], \sigma_A) \leq \| (\langle A^d \rangle \bigcirc □ \varphi) \|
\]

\[
\Rightarrow \forall q' \in \text{out}(\lambda[i], \sigma_A),
\]

\[
\exists d'-\text{strategy } F_{A,q'}.
\]

\[
\forall \lambda' \in \text{out}(q', F_{A,q'}), \forall j \geq 1: \lambda'[j] \in \| \varphi \|
\]

Then, we define for every \( \lambda \in \lambda^i \) that \( F_A(\lambda) = \sigma_A \) and \( \Lambda^{i+1} = \{ \lambda q' | \lambda \in \lambda^i \land q' \in \text{out}(\lambda[i], \sigma_A) \} \). Obviously, we have:

\[
\forall \lambda q' \in \lambda^{i+1}: \text{cost}(\lambda, F_A) = \text{cost}(\lambda, F_A) + \text{cost}(\lambda[i], \sigma_A)
\]

\[
\leq (0^\infty \leftarrow b) + d = d
\]

\[
\land
\]

\[
q' \in \| (\langle A^d \rangle \bigcirc □ \varphi) \|
\]

* If \( \lambda[i] \in \| (\langle A^0 \rangle \bigcirc b) \| \) \( \bigcirc (\langle Z \rangle) \), we have:

\[
\lambda[i] \in \| (\langle A^0 \rangle \bigcirc \bigcirc (\langle Z \rangle) \Rightarrow \exists \sigma_{A,\lambda} \in D_A(\lambda[i]): \text{cost}(\lambda[i], \sigma_{A,\lambda}) \leq 0^\infty \leftarrow b \quad \text{and}
\]

\[
\text{out}(\lambda[i], \sigma_{A,\lambda}) \subseteq Z
\]

Then, we define for every \( \lambda \in \lambda^i \) that \( F_A(\lambda) = \sigma_{A,\lambda} \) and \( \Lambda^{i+1} = \{ \lambda q' | \lambda \in \lambda^i \land q' \in \text{out}(\lambda[i], \sigma_{A,\lambda}) \} \). Obviously, we have:

\[
\forall \lambda q' \in \lambda^{i+1} \Rightarrow \text{cost}(\lambda q', F_A) = \text{cost}(\lambda, F_A) + \text{cost}(\lambda[i], \sigma_{A,\lambda})
\]

\[
\land
\]

\[
q' \in \text{out}(\lambda[i], \sigma_{A,\lambda})
\]

\[
\Rightarrow \text{cost}(\lambda q', F_A) \leq 0^\infty \leftarrow b \land
\]

\[
q' \in Z \quad \text{as } \text{out}(\lambda[i], \sigma_{A,\lambda}) \subseteq Z
\]
Given the above construction of $F_A$, we have that 
\[ \forall \lambda \in \text{out}(q, F_A) \text{ and } i \geq 1 \Rightarrow \lambda \in \Lambda^i \]
\[ \Rightarrow \exists (d, d') \in \text{split}(b) : \]
\[ \text{cost}(\lambda[i], F_A) \leq d \wedge \]
\[ \lambda[i] \in \| (A^d) \Box \varphi \| \]
or
\[ \text{cost}(\lambda[i], F_A) \leq \bar{\beta} \rightarrow b \wedge \lambda[i] \in Z \]
\[ \Rightarrow \text{cost}(\lambda[i], F_A) \leq b \wedge \lambda[i] \in \| \varphi \| \]

In other words, $q \in \| (A^b) \Box \varphi \|$, i.e., $q \in Y$.

Therefore, $Z \subseteq Y$; hence, $Y$ is the greatest post-fixed point of $f(X)$, hence also the greatest fixed point of $f(X)$. ■

3.2 Soundness of RB-ATL

First, we prove that the axioms of RB-ATL are valid.

(⊥) is valid because there is no $b$-strategy $F_A$ such that for all $\lambda \in \text{out}(q, F_A)$, $\lambda[1]$ makes $\bot$ true.

(⊤) is valid because $A$ has a $\bar{\beta}$-strategy $F_A$ such that for all $\lambda \in \text{out}(q, F_A)$, $\lambda[1]$ makes $\top$ true.

(B) is valid because if there is a $b$-strategy $F_A$ such that for all $\lambda \in \text{out}(q, F_A)$, $\lambda[1]$ makes $\varphi$ true, then the same $F_A$ is also a $d$-strategy which has the same property.

(S) is valid because if there exists a strategy $F_{A_1}$ to enforce $\varphi$ and a strategy $F_{A_2}$ to enforce $\psi$, then there exists a joint strategy $F_{A_1} \cup A_2$ (with the same moves for $A_1$ and $A_2$ as $F_{A_1}$ and $F_{A_2}$, respectively) to enforce both $\varphi$ and $\psi$.

(SN) is valid because if there exists a $b$-strategy $F_A$ to enforce $\varphi$, and, for all strategies of $N$, $\psi$ is inevitable, then $\varphi \wedge \psi$ can be enforced in by $F_A$.

(SN+) is valid because if there exists a $b$-strategy $F_N$ to enforce $\varphi$, and, for all strategies of $N$ which cost at most $b$, $\psi$ is inevitable, then $\varphi \wedge \psi$ can be enforced in by $F_N$.

(SN-) is valid because if, for all strategies of $N$, $\varphi$ and $\psi$ are inevitable, then so is $\varphi \wedge \psi$.

(FP2) is valid by Lemma 1(1) and (FPU) by Lemma 1(2).

Then, we prove that the inference rules preserve validity (the proof for (MP) is standard, hence it is omitted):

((A^b) □-Monotonicity), ((A^b) □-Monotonicity), and ((A^b) U-Monotonicity) clearly preserve validity, since if $\| \varphi \| \subseteq \| \psi \|$ and an outcome in $\| \psi \|$ can be enforced, then an outcome in $\| \psi \|$ can also be enforced by the same strategy.

((N^b) □-Necessitation) is valid since if $\varphi$ is logically true, then it is inevitable in perpetuity.

((A^b) □-Induction) and ((A^b) U-Induction) preserve validity by Lemma 1.

3.3 Completeness of RB-ATL

The proof of completeness is based on [13]. We construct a satisfying model for a formula $\varphi_0$ which is consistent with the axiomatic system for RB-ATL.
In the proof, we assume when convenient that all formulas are in negation normal form of RB-ATL. The syntax of negation normal form RB-ATL is as follows:

\[
\phi ::= p | \neg p | p \lor q | p \land q | (\langle A \rangle^{b}) \circ \phi | (\langle A \rangle^{b}) \Box \phi | (\langle A \rangle^{b}) \Diamond \phi | (\langle A \rangle^{b}) \vartriangleleft \Diamond \phi
\]

where \( A \) is a non-empty coalition and \( b \in \mathbb{B} \). Given a normal form formula \( \phi \) of RB-ATL, we denote by \( \overline{\phi} \) the normal form negation of \( \phi \). Given an RB-CGS \( S \) and a state \( q \), the semantics of normal form RB-ATL is the same as RB-ATL, except for formulas \( \overline{(\langle A \rangle^{b})} \circ \phi \), \( \overline{(\langle A \rangle^{b})} \Box \phi \) and \( \overline{(\langle A \rangle^{b})} \Diamond \phi \) which are defined as follows:

- \( S, q \models \overline{(\langle A \rangle^{b})} \circ \phi \) iff for every \( b \)-strategy \( F_{A} \), there exists \( \lambda \in \operatorname{out}(q, F_{A}) \) such that \( S, \lambda[1] \models \neg \phi \) iff \( \forall \sigma_{A} \in D_{A}(q) : \text{cost}(q, \sigma_{A}) \leq b \rightarrow \exists q' \in \operatorname{out}(q, \sigma) : S, q' \models \neg \phi. \)
- \( S, q \models \overline{(\langle A \rangle^{b})} \Box \phi \) iff for every \( b \)-strategy \( F_{A} \), there exists \( \lambda \in \operatorname{out}(q, F_{A}) \) and \( i \geq 1 \) such that \( S, \lambda[i] \models \neg \phi \).
- \( S, q \models \overline{(\langle A \rangle^{b})} \vartriangleleft \Diamond \phi \) iff for every \( b \)-strategy \( F_{A} \), there exists \( \lambda \in \operatorname{out}(q, F_{A}) \) and if there exists \( i \geq 1 \) with \( S, \lambda[i] \models \psi \) then there is \( j \in \{1, \ldots, i-1\} \) where \( S, \lambda[j] \models \neg \phi \).

The model is constructed in a way very similar to the construction in [13]. It is assembled from finite trees where nodes are labelled by sets of formulas. First, we define the set of formulas used in the labelling.

**Definition 6**

The closure \( cl(\varphi_{0}) \) is the smallest set of formulas satisfying the following closure conditions:

- all sub-formulas of \( \varphi_{0} \) including \( \varphi_{0} \) itself are in \( cl(\varphi_{0}) \);  
- if \( (\langle A \rangle^{d}) \Box \phi \) is in \( cl(\varphi_{0}) \), then so are \( (\langle A \rangle^{d}) \circ (\langle A \rangle^{d}) \Box \phi \) for all \( (d, d') \in \text{split}(b) \) and also \( (\langle A \rangle^{d_{\leq b}^{d}}) \circ (\langle A \rangle^{b}) \Box \phi \);  
- if \( (\langle A \rangle^{b}) \vartriangleleft \Diamond \phi \) is in \( cl(\varphi_{0}) \), then so are \( (\langle A \rangle^{d}) \circ (\langle A \rangle^{d}) \vartriangleleft \Diamond \phi \) for all \( (d, d') \in \text{split}(b) \) and also \( (\langle A \rangle^{d_{\leq b}^{d}}) \circ (\langle A \rangle^{b}) \vartriangleleft \Diamond \phi \);  
- if \( \psi \) is in \( cl(\varphi_{0}) \), then so is \( \neg \psi \); and  
- \( cl(\varphi_{0}) \) is also closed under finite positive Boolean operators (\( \lor \) and \( \land \)) up to tautology equivalence.

Note that \( cl(\varphi_{0}) \) is finite. Let \( \Gamma \) be the set of maximal consistent subsets of \( cl(\varphi_{0}) \). We define trees \( (T, V, C) \overline{\Gamma} \) over \( \Gamma \) in a similar way as [13] where

- \( T \subseteq (\mathbb{N}^{\ast})^{\ast} \) is the set of nodes;  
- \( V : T \rightarrow \Gamma \) is a labelling function that assigns to each node a consistent set; and  
- \( C : T \times N \times N \rightarrow \mathbb{N}^{\ast} \) is a (partial) cost function that assigns a cost to each action available at a node.

Intuitively, nodes in a tree are identified with finite words corresponding to the sequence of joint actions by the grand coalition which leads to that node. The root is the empty word \( e \) and each node \( t \) corresponds to a finite computation the last state of which is \( t \). An interior node of the tree is a node but not a leaf. A formula is in \( V(t) \) intuitively means that the formula is true in \( t \). Finally, the cost of an action \( j \) of an agent \( i \) at a node \( t \) is given by \( C(t, i, j) \). As in [13], the construction proceeds in three stages. The first stage is producing locally consistent trees, namely trees where the labelling satisfies...
conditions on successor nodes which makes it possible to prove a truth lemma for the next
step modalities. The second stage is proving the existence of trees which realize eventualities (essentially,
make the labelling consistent with the truth conditions for the □ and U modalities). Finally, the finite
trees realizing eventualities are combined into one infinite tree model.

**Definition 7**
A tree (T, V, C) is locally consistent iff for any interior node t ∈ T:

1. If ⟨⟨Abi⟩⟩⃝ϕ in V(t), then there is a move σA such that cost(t, σA) ≤ b and for any t′ ∈ out(t, σA)
   we have ϕ ∈ V(t′); and
2. If ¬⟨⟨Ab⟩⟩⃝ϕ in V(t), then for any move σA with cost(t, σA) ≤ b, there exists t′ ∈ out(t, σA)
   where ¬ϕ ∈ V(t′).

Two following lemmas are used as a crucial step in the local consistency proof.

**Lemma 2**
Let \( \Phi = \{ ⟨⟨Ab1⟩⟩⃝ϕ1, ..., ⟨⟨Abk⟩⟩⃝ϕk, ¬⟨⟨N\bar{∞}⟩⟩⃝χ1, ..., ¬⟨⟨N\bar{∞}⟩⟩⃝χm, ¬⟨⟨Ab⟩⟩⃝ψ \} \) be a consistent
set of formulas in which:

- all \( Ai \) are both non-empty and pair-wise disjoint
- \( \bigcup_i A_i \subseteq A \)
- \( \sum_i b_i ≤ b \)

Then, \( \Psi = \{ ϕ1, ..., ϕ_k, ¬χ1, ..., ¬χ_m, ¬ψ \} \) is also consistent.

**Proof.** When \( k = 0 \) (or \( m = 0 \)), we can always add the axiom ⟨⟨Ab′⟩⟩⃝T (or ¬⟨⟨N\bar{∞}⟩⟩⃝⊥) into \( \Phi \).
Hence, it is sufficient to prove this lemma with \( k > 0 \) and \( m > 0 \).

Let \( A′ = \bigcup_i A_i, b′ = \sum b_i \) and \( ϕ = \bigwedge_i ϕ_i \) and \( χ = \bigwedge_j ¬χ_j \).
Assume to the contrary that \( \Psi \) is inconsistent, we have:

1. \( \vdash \bigwedge_i ⟨⟨Ab'_i⟩⟩⃝ϕ_i → ⟨⟨Ab′⟩⟩⃝ϕ \)
   by (S), (B), \( A′ \subseteq A \) and \( b′ ≤ b \)
2. \( \vdash \bigwedge_j ¬⟨⟨N\bar{∞}⟩⟩⃝χ_j → ¬⟨⟨N\bar{∞}⟩⟩⃝¬χ \)
   by (SN−)
3. \( \vdash ⟨⟨Ab′⟩⟩⃝ϕ \land ¬⟨⟨N\bar{∞}⟩⟩⃝¬χ → ⟨⟨Ab′⟩⟩⃝(ϕ \land χ) \)
   by (SN)
4. \( \vdash \bigwedge_i ⟨⟨Ab'_i⟩⟩⃝ϕ_i \land \bigwedge_j ¬⟨⟨N\bar{∞}⟩⟩⃝χ_j → ⟨⟨Ab′⟩⟩⃝(ϕ \land χ) \)
   by (1), (2) and (3)
5. \( ϕ \land χ → ψ \)
   as \( \Psi \) is inconsistent
(6) \( \vdash \langle \langle A^b \rangle \rangle \circ (\phi \land \chi) \rightarrow \langle \langle A^b \rangle \rangle \circ \psi \)

by (5) and \( \langle \langle A^b \rangle \rangle \circ \)-monotonicity

(7) \( \vdash \bigwedge_i \langle \langle A^b_i \rangle \rangle \circ \phi_i \land \bigwedge_j \neg \langle \langle N^\infty \rangle \rangle \circ \chi_j \rightarrow \langle \langle A^b \rangle \rangle \circ \psi \)

by (4) and (6)

Therefore, \( \Phi \cup \{ \langle \langle A^b \rangle \rangle \circ \psi \} \) is consistent, which is a contradiction.

Similarly, we have the following lemma:

LEMMA 3

Let \( \Phi = \{ \langle \langle A_1^{b_1} \rangle \rangle \circ \phi_1, \ldots, \langle \langle A_k^{b_k} \rangle \rangle \circ \phi_k, \neg \langle \langle N^d \rangle \rangle \circ \chi_1, \ldots, \neg \langle \langle N^d \rangle \rangle \circ \chi_m \} \) be a consistent set of formulas in which:

- The \( A_i \)'s are both non-empty and pair-wise disjoint
- \( \sum_i b_i \leq d_j \) for all \( j \)

Then, \( \Psi = \{ \psi_1, \ldots, \psi_k, \neg \chi_1, \ldots, \neg \chi_m \} \) is also consistent.

PROOF. When \( k = 0 \) (or \( m = 0 \)), we can always add the axiom \( \langle \langle N^0 \rangle \rangle \circ \top \) (or \( \neg \langle \langle N^\infty \rangle \rangle \circ \bot \)) into \( \Phi \). Hence, it is sufficient to prove this lemma with \( k > 0 \) and \( m > 0 \).

Let \( A = \bigcup_i A_i, b = \sum_i b_i \) and \( \psi = \bigwedge_i \phi_i \) and \( \chi = \bigwedge_j \neg \chi_j \).

Assume to the contrary that \( \Psi \) is inconsistent, we have:

(1) \( \vdash \bigwedge_i \langle \langle A_i^{b_i} \rangle \rangle \circ \phi_i \rightarrow \langle \langle A^b \rangle \rangle \circ \psi \)

by (S)

(2) \( \vdash \neg \langle \langle N^d \rangle \rangle \circ \chi_j \rightarrow \neg \langle \langle A^b \rangle \rangle \circ \chi_j \)

by B and \( b \leq d_j \)

(3) \( \vdash \bigwedge_j \neg \langle \langle N^b \rangle \rangle \circ \chi_j \rightarrow \neg \langle \langle A^b \rangle \rangle \circ \neg \chi \)

by (S\text{\textendash}N\textendash)

(4) \( \vdash \langle \langle A^b \rangle \rangle \circ \psi \land \neg \langle \langle N^b \rangle \rangle \circ \neg \chi \rightarrow \langle \langle A^b \rangle \rangle \circ (\phi \land \chi) \)

by (S\text{\textendash}N\textendash)

(5) \( \vdash \bigwedge_i \langle \langle A_i^{b_i} \rangle \rangle \circ \phi_i \land \bigwedge_j \neg \langle \langle N^b \rangle \rangle \circ \chi_j \rightarrow \langle \langle A^b \rangle \rangle \circ (\phi \land \chi) \)

by (1), (2), (3) and (4)

(6) \( \vdash \phi \land \chi \rightarrow \bot \)

as \( \Psi \) is inconsistent
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\[ (7) \vdash (\langle A^b \rangle) (\varphi \land \chi) \rightarrow (\langle A^b \rangle) \circ \bot \]

by (6) and \( (\langle A^b \rangle) \circ \text{-monotonicity} \)

\[ (8) \vdash \bigwedge_i (\langle A^b_i \rangle) \circ \varphi_i \land \bigwedge_j \neg (\langle N^{b_j} \rangle) \circ \chi_j \rightarrow (\langle A^b \rangle) \circ \bot \]

by (5) and (7)

Therefore, \( \Phi \cup \{ (\langle A^b \rangle) \circ \bot \} \) is consistent, which is a contradiction \( \blacksquare \)

**Lemma 4**

Let \( \Phi \) be a finite consistent set of formulas. Let \( \Phi_\circ \) be the subset of \( \Phi \) which contains all formulas of the form \( (\langle A^b \rangle) \circ \varphi \) or their negations. Let \( k \in \mathbb{N} \) be such that \( |\Phi_\circ| < k \), then there is a locally consistent tree \( (T, V, C) \) of height one where \( T = (\epsilon) \cup \{1, \ldots, k\}^n \) and \( V(\epsilon) = \Phi \).

**Proof.** Denote \( T' = \{1, \ldots, k\}^n \); hence \( T = (\epsilon) \cup T' \) where we denote by \( \cup \) the disjoint union operator. Furthermore, we assume that

\[ \Phi_\circ = \Phi^+_\circ \cup \Phi^-_\circ \cup \Phi^N_\circ \]

where

\[ \Phi^+_\circ = \{(\langle A^b_1 \rangle) \circ \varphi_1, \ldots, (\langle A^b_m \rangle) \circ \varphi_m \}, \]

\[ \Phi^-_\circ = \{\neg (\langle B^{a_i}_1 \rangle) \circ \varphi_1, \ldots, \neg (\langle B^{a_i}_n \rangle) \circ \varphi_l \} \quad \text{s.t. } \forall i : B_i \neq N \]

and

\[ \Phi^N_\circ = \{\neg (\langle N^{c_i} \rangle) \circ \chi_1, \ldots, \neg (\langle N^{c_k} \rangle) \circ \chi_k \} \]

First, let \( f \in \mathbb{N} \) be the maximal number which occurs in \( e_1, \ldots, e_5 \); we define \( f + 1 = (f + 1)^n \). It is straightforward that for all \( e \in \{e_1, \ldots, e_5\} \), if \( e \neq \infty \) then \( f + 1 \leq e \). We define a function \( \text{deinf} : \mathbb{B}^\infty \to \mathbb{B} \) which removes infinity from a bound as follows: \( \text{deinf}(b) = b' \) where for all \( i = 1, \ldots, r \)

\[ b'_i = \begin{cases} b_i & \text{if } b_i \neq \infty \\ f + 1 & \text{otherwise} \end{cases} \]

It is also straightforward that for all \( e \in \{e_1, \ldots, e_5\} \) and \( b \in \{b_1, \ldots, b_m\} \), if \( e \neq \infty \) and \( b \notin \mathbb{N} \) (i.e. \( b \) contains some \( \infty \)) then \( \text{deinf}(b) \leq e \).

Let us construct a tree with a root labelled by \( \Phi \) and \( k^n \) children denoted by \( t = (a_1, \ldots, a_n) \in \{1, \ldots, k\}^n \). Intuitively, we allow each agent to perform \( k \) different actions where the special action \( k \) always costs \( 0 \). For convenience, we denote the action of agent \( i \) in \( t \) by \( t_i = a_i \) and the joint move by a coalition \( A \) in \( t \) by \( t_A = (t_i)_{i \in A} \). In the following, we define the labelling function \( V(t) \) for each leaf \( t \) and the cost function \( C(e, i, a) \) for each agent \( i \in N \) and action \( a \in \{1, \ldots, k\} \):

(a) For each \( (\langle A^b_p \rangle) \circ \varphi_p \in \Phi^+_\circ \) where \( A_p \neq \emptyset, \varphi_p \) is added to \( V(t) \) for all \( t \) such that \( \forall i \in A_p : t_i = p \).

Let \( \min_{A_p} \) be the smallest number in \( A_p \), we assign the cost of action \( p \) performed by \( \min_{A_p} \)
to be \textit{deinf}(b_p), i.e. \( C(\epsilon, \min A_p, p) = \text{deinf}(b_p) \). For other agents \( i \) in \( A_p \setminus \{ \min A_p \} \), we assign \( C(\epsilon,i,p)=0 \). For other unassigned-cost actions, their costs are assigned as follows:

\[
C(\epsilon,i,p)=\begin{cases} 0 & \text{if } p \leq m \text{ or } p = k \\ \frac{1}{f+1} & \text{if } m < p < k \\
\end{cases}
\]

We define \( C(t,A) = \sum_{i \in A} C(\epsilon, i, t_i) \) as the cost of the joint action by the coalition \( A \) and \( C(t) = C(t, N) \) as the cost of the joint action by the grand coalition.

(b) For each \( \neg(\langle N^r \rangle) \bigwedge \chi_p \in \Phi_{\nabla} \), \( \neg \chi_p \) is added to \( V(t) \) whenever \( C(t) \leq e_p \).

(c) Finally, we will add at most one formula from \( \Phi_{\nabla} \) to \( V(t) \). Let

\[
\Phi_{\nabla}(t) = \{ \neg(\langle B^d \rangle) \bigwedge \psi \in \Phi_{\nabla} \mid C(t,B) \leq d \}
\]

where \( 1 \leq i_t < i_2 < \ldots < i_l \leq l \). Let \( I = \{ i' \mid i' \in \{ m+1, \ldots, m+l \} \} \) and \( j = \sum_{i \in I}(i_t - m - 1) \) mod \( l_t + 1 \). Consider \( \neg(\langle B^d_{j_t} \rangle) \bigwedge \psi_{i_t} \): if \( N \setminus B_j \subseteq I \), then \( \neg \psi_{i_t} \) is added into \( V(t) \).

We now prove that the constructed tree \( (T,V,C) \) is locally consistent. First, we show that all labels are consistent. It is obvious that \( V(\epsilon) = \Phi \) is consistent. For each child \( t \in T' \), since at most one formula \( \neg \psi_p \) such that \( \neg(\langle B^d_p \rangle) \bigwedge \psi_p \in \Phi_{\nabla} \) is added into \( V(t) \), we consider the following two cases:

\textbf{Case} \( \forall p \in \{ 1, \ldots, l \} : \neg \psi_p \notin V(t) \) : Let us assume that

\[
V(t) = \{ \psi_{i_1}, \ldots, \psi_{i_m} \} \cup \{ \neg \chi_{j_1}, \ldots, \neg \chi_{j_h} \}
\]

where \( 1 \leq i_1 \leq \ldots \leq i_m \leq m \) and \( 1 \leq j_1 \leq \ldots \leq j_h \leq m \). Then, we have:

(a) \( \Rightarrow \forall p \in \{ i_1, \ldots, i_m \} : i_t = p \Rightarrow \forall p, p' \in \{ i_1, \ldots, i_m \} : p \neq p' \Rightarrow A_p \cap A_{p'} = \emptyset \)

and if \( \forall p \in \{ i_1, \ldots, i_m \} : b_p \in \mathbb{N}^r \), then

(a) \( \Rightarrow \forall p \in \{ i_1, \ldots, i_m \} : \text{deinf}(b_p) = b_p = C(t,A_p) \)

\[
\Rightarrow \sum_{p \in \{ i_1, \ldots, i_m \}} b_p = \sum_{p \in \{ i_1, \ldots, i_m \}} C(t,A_p) \leq C(t)
\]

\( \Rightarrow \forall j \in \{ j_1, \ldots, j_h \} \sum_{p \in \{ i_1, \ldots, i_m \}} b_p \leq e_j \) by (b)
ATL with resource bounds

otherwise, if $\exists p \in \{i_1, \ldots, i_m\} : b_p \not\in N'$, then

(a) $\Rightarrow \forall p \in \{i_1, \ldots, i_m\} : \text{deinf}(b_p) = C(t, A_p)$
   $\Rightarrow \forall j \in \{1, \ldots, h\} : e_j \neq \infty \rightarrow$
   $\forall p \in \{i_1, \ldots, i_m\}$
   $\sum \text{deinf}(b_p) = \sum C(t, A_p) \leq e_j$
   $\Rightarrow \forall j \in \{1, \ldots, h\} : \sum_{p \in \{i_1, \ldots, i_m\}} b_p \leq e_j$

Therefore, by Lemma 3, $V(t)$ is consistent.

Case $\exists! q \in \{1, \ldots, l\} : \neg \psi_q \in V(t)$ :

Let us assume that

$$V(t) = \{\psi_{i_1}, \ldots, \psi_{i_m}\} \cup \neg \psi_q \cup \neg \chi_{j_1} \ldots \neg \chi_{j_h}$$

where $1 \leq i_1 \leq \ldots \leq i_m \leq m$ and $1 \leq j_1 \leq \ldots \leq j_h \leq m$. Recall that:

$$\Phi_\bigcirc(t) = \neg (\langle B_{q_1}^{d_1} \rangle \psi_{q_1} \ldots \neg (\langle B_{q_h}^{d_h} \rangle \psi_{q_h})$$

hence, $q \in \{q_1, \ldots, q_l\}$, and also

$$I = \{i' | t_{i'} \in \{m+1, \ldots, m+l\}\}$$

hence, $\forall i \in \{i_1, \ldots, i_m\} : I \cap A_i = \emptyset$ since $t_{i'} = i$ for all $i' \in A_i$.

Similar to the previous case, we have:

(a) $\Rightarrow \forall p \in \{i_1, \ldots, i_m\}, \forall i \in A_p : t_i = p$
   $\Rightarrow \forall p, p' \in \{i_1, \ldots, i_m\} : p \neq p' \rightarrow A_p \cap A_{p'} = \emptyset$

Then, we have:

(c) $\Rightarrow N \setminus B_q \subseteq I$
   $\Rightarrow I \neq \emptyset$ as $B_q \neq N$
   $\Rightarrow C(t) \geq C(t, I) \geq \overline{f+1}$ by (a)
   $\Rightarrow \forall j \in \{1, \ldots, h\} : e_j \neq \infty \rightarrow C(t) \leq e_j$
   as $\overline{f+1} \leq e_j$
   $\Rightarrow \forall j \in \{j_1, \ldots, j_h\} : e_j = \infty$ by (b)
and

\[(c) \Rightarrow N \setminus B_q \subseteq I\]
\[\Rightarrow B_q \supseteq N \setminus I\]
\[\Rightarrow \forall i \in \{i_1, \ldots, i_m\}: B_q \supseteq A_i \quad \text{as } I \cap A_i = \emptyset\]
\[\Rightarrow B_q \supseteq A_{i_1} \cup \ldots \cup A_{i_m}\]
\[\Rightarrow C(t, B_q) \geq C(t, A_{i_1} \cup \ldots \cup A_{i_m})\]
\[\Rightarrow C(t, B_q) \geq C(t, A_{i_1}) + \ldots + C(t, A_{i_m})\]
\[\Rightarrow d_q \geq \sum_{i = \{i_1, \ldots, i_m\}} b_i\]

Therefore, by Lemma 2, \(V(t)\) is consistent.

Let us now prove that \((T, V, C)\) satisfies the two local consistency conditions of Definition 7.

(1.) Assume that \(\langle\langle A_p^{b_p} \rangle\rangle \circ \psi_p \in V(\epsilon)\). Consider the joint action \(\sigma\) for \(A_p\) such that \(\sigma_i = p\) for all \(i \in A_p\). We have:

\[\text{out}(\epsilon, \sigma) = \{t \in T' | \forall i \in A_p : t_i = p\}\]

and

\[\langle\langle A_p^{b_p} \rangle\rangle \circ \psi_p \in V(\epsilon) \Rightarrow \langle\langle A_p^{b_p} \rangle\rangle \circ \psi_p \in \Phi^+\]
\[\Rightarrow \forall t \in T': (\forall i \in A_p : t_i = p) \rightarrow \psi_p \in V(t) \quad \text{by (a)}\]
\[\Rightarrow \forall t \in \text{out}(\epsilon, \sigma) : \psi_p \in V(t)\]

(2.) If \(\neg \langle\langle N^{\psi} \rangle\rangle \circ \chi_p \in V(\epsilon)\), let us consider the joint action \(t\) such that \(C(t) \leq e_p\). Obviously, \(\text{out}(\epsilon, t) = \{t\}\). Hence, by (b), \(\sim \chi_p \in V(t)\).

If \(\neg \langle\langle B_p^{d_p} \rangle\rangle \circ \psi_p \in V(\epsilon)\) where \(B_p \neq N\), let \(\sigma\) be an arbitrary joint move for the coalition \(B_p\) such that \(C(t, \sigma) \leq d_p\). We will determine an outcome \(t \in \text{out}(\epsilon, \sigma)\), such that \(\sim \psi_p \in V(t)\).

As \(t \in \text{out}(\epsilon, \sigma)\), \(t_i = \sigma_i\) for all \(i \in B_p\); it remains to determine \(t_i\) for \(i \notin B_p\).

Let \(t' \in T'\) such that

\[t_i' = \begin{cases} \sigma_i & \text{if } i \in B_p \\ m + 1 & \text{otherwise} \end{cases}\]

and

\[\Phi^-(\sigma) = \Phi^-(t')\]
\[= \{\neg \langle\langle B_i^{d_i} \rangle\rangle \circ \psi_{t_i}, \ldots, \neg \langle\langle B_{i_{l_p}}^{d_{i_{l_p}}} \rangle\rangle \circ \psi_{t_{i_{l_p}}} \} \]
\[I_\sigma = \{i \in B_p | \sigma_i \in [m + 1, \ldots, m + l] \}\]
\[j_\sigma = \sum_{i \in I_\sigma} (\sigma_i - m - 1)\]

Since \(C(t', B_p) \leq d_p\), \(\neg \langle\langle B_p^{d_p} \rangle\rangle \circ \psi_p \in \Phi^-(\sigma)\). Let \(p = j_\sigma\) for some \(j_\sigma \in [1, \ldots, l_\sigma]\).
Let $i$ be an arbitrary agent in $N \setminus B_p \neq \emptyset$. We define:

$$t_i = \begin{cases} m + (j_i - j_{\sigma} - 1) \mod l_{\sigma} + 1 & \text{if } i = \iota \\ m + 1 & \text{if } i \in N \setminus B_p \setminus \{\iota\} \end{cases}$$

Let us prove that $\sim \psi_p \in V(t)$.

We have:

$$\forall i \in B_p : t_i = \sigma_i \Rightarrow C(\epsilon, i, t_i) = C(\epsilon, i, t'_i)$$

and

$$\forall i \in N \setminus B_p : t_i \in \{m + 1, \ldots, m + l\} \Rightarrow C(\epsilon, i, t_i) = \sigma + 1 \quad \text{by (a)}$$

$$\Rightarrow C(\epsilon, i, t_i) = C(\epsilon, i, t'_i)$$

Therefore, $\Phi_{\sigma}(t) = \Phi_{\sigma}(t') = \Phi_{\sigma}(\sigma)$; hence, $l = l_{\sigma}$. Then,

$$I = \{i | t_i \in \{m + 1, \ldots, m + l\}\}$$

$$= I_{\sigma} \cup (N \setminus B_p)$$

$$j = \sum_{i \in I} (t_i - m - 1) \mod l_{\sigma} + 1$$

$$= (\sum_{i \in I_{\sigma}} (t_i - m - 1)) +$$

$$\quad (t_i - m - 1) +$$

$$\quad \sum_{i \in N \setminus B_p \setminus \{\iota\}} (t_i - m - 1) \mod l_{\sigma} + 1$$

$$= (j_{\sigma} + (j_{\sigma} - j_{\sigma} - 1) \mod l_{\sigma}) \mod l_{\sigma} + 1$$

$$= (j_{\sigma} - 1) \mod l_{\sigma} + 1$$

$$= j_{\sigma}$$

Recall that $j_{\sigma} = p$. Then, we have:

$$N \setminus B_p \subseteq N \setminus B_p \cup I_{\sigma} = I$$

Therefore, according to (c), $\sim \psi_p \in V(t)$. 

The next stage of the proof is to consider what conditions on tree labelling we need to be able to prove the truth lemma for other temporal modalities. The definition of what it means to ‘realize’ formulas of the form $\langle\langle A^b \rangle\rangle \psi U \psi$, $\neg \langle\langle A^b \rangle\rangle \Box \phi$, $\langle\langle A^b \rangle\rangle \Box \phi$, $\neg \langle\langle A^b \rangle\rangle \phi U \psi$ is similar to the one in [13] (essentially the truth conditions for the formulas with ‘satisfied’ replaced by ‘in the labelling of’).

The following lemma and its proof are similar to the correspondings in [13], but for formulas of RB-ATL.
**Lemma 5**
For any subset $Y \subseteq \Gamma_1$, there is a formula $\chi_Y \in \text{cl}(\varphi_0)$, called the characteristic formula of $Y$, such that for every $y \in \Gamma_1$, $\chi_Y \in y$ iff $y \in Y$.

**Proof.** For any maximally consistent subset $y$ of $\text{cl}(\varphi_0)$, we define:

$$X_{\{y\}} = \bigwedge_{\varphi \in y} \varphi \quad \text{and} \quad X_Y = \bigvee_{y \in Y} X_{\{y\}}$$

First, since $\text{cl}(\varphi_0)$ is closed under finite positive boolean operators, $X_{\{y\}} \in y$.

$(\Rightarrow)$: Assume that $\chi_Y \in y$. Then, we have:

$$\chi_Y \in y \Rightarrow \vdash X_{\{y\}} \rightarrow \chi_Y \quad \text{a PL tautology}$$

$$\Rightarrow \chi_Y \in y \quad \text{as } X_{\{y\}} \in y \text{ and } y \text{ is maximal}$$

$(\Rightarrow)$: Assume that $y \not\in Y$. Then, for any $y' \in Y$, we have:

$$\exists \theta \in \text{cl}(\varphi_0): \theta \in y' \text{ and } \sim \theta \in y$$

$$\Rightarrow \chi_{\{y'\}} \wedge \sim \theta \text{ is inconsistent}$$

$$\Rightarrow \chi_{\{y'\}} \not\in y$$

Therefore, $\chi_Y \not\in y$.

In what follows, $\Psi_\boxempty$ is the set of formulas of the form $\langle\langle \text{Ab} \rangle\rangle \varphi$ or $\neg \langle\langle \text{Ab} \rangle\rangle \varphi$ from $\text{cl}(\varphi_0)$.

**Lemma 6**
Given $\langle\langle \text{Ab} \rangle\rangle \psi \cup \eta$ and $x \in \Gamma_1$, there is finite tree $(T, V, C)$ over $\Gamma_1$ such that:

- every interior node of $(T, V, C)$ has $k^n$ children where $k = |\Psi_\boxempty| + 1$,
- $(T, V, C)$ is locally consistent,
- $V(\epsilon) = x$, and
- if $\langle\langle \text{Ab} \rangle\rangle \psi \cup \eta \in x$ then $(T, V, C)$ realizes $\langle\langle \text{Ab} \rangle\rangle \psi \cup \eta$ from $\epsilon$

**Proof.** The proof is similar to the corresponding proof in [13], but also uses induction on the bound $b$.

Let $Z \subseteq \Gamma_1$ such that, for any $x \in Z$, there is a finite tree obeying all conditions of Lemma 6. We shall prove the lemma by showing that $Z = \Gamma_1$. In the following, assume that $x \in \Gamma_1$.

- If $\langle\langle \text{Ab} \rangle\rangle \psi \cup \eta \not\in x$, let us construct a simple tree $(T, V, C)$ where $T = \{\epsilon\}$ and $V(\epsilon) = x$. Since $(T, V, C)$ satisfies all conditions of Lemma 6, $x \in Z$.
- If $\langle\langle \text{Ab} \rangle\rangle \psi \cup \eta \in x$, we first show that $\eta = (\psi \wedge (\langle\langle \text{Ab} \rangle\rangle \psi)$ is a theorem. Then, we have that:

$$(1) \vdash (\psi \wedge (\langle\langle \text{Ab} \rangle\rangle \psi) \rightarrow \chi_Z$$

$$(2) \vdash \langle\langle \text{Ab} \rangle\rangle \psi \rightarrow \chi_Z \quad \text{by } \langle\langle \text{Ab} \rangle\rangle \text{-Induction}$$

Therefore,

$$\langle\langle \text{Ab} \rangle\rangle \psi \in x \Rightarrow \chi_Z \in x \quad \text{as } x \text{ is maximal}$$

$$\Rightarrow x \in Z \quad \text{by Lemma 5}$$
However, it remains to prove that $\eta$ is a theorem. This is done by showing that $\eta$ belongs to any maximal consistent set $q$ of RB-ATL in three cases:

- If $\langle\langle A^b \rangle\rangle \psi \psi \not\in q$, let us construct a simple tree $(T, V, C)$ where $T = \{e\}$ and $V(e) = q \cap cl(\phi_0)$. Since $(T, V, C)$ satisfies all conditions of the lemma, $q \cap cl(\phi_0) \in Z$. Then, we have:

$$q \cap cl(\phi_0) \in Z \Rightarrow \chi(\phi_0) \in q \cap cl(\phi_0) \quad \text{by Lemma 5}$$

$$\Rightarrow \chi \in q \quad \text{as } q \text{ is maximal}$$

- If $\psi \lor (\psi \land \langle\langle (A^b) \rangle\rangle) \not\in q$, then it is straightforward that $\eta \in q$ since $q$ is maximal.

- If $\langle\langle (A^b) \rangle\rangle \psi \not\in q$ and $\psi \lor (\psi \land \langle\langle (A^b) \rangle\rangle) \not\in q$, we prove that $\eta \in q$ by induction on $b$.

**Base case:** Assume that $b = 0^\infty$. As $\psi \lor (\psi \land \langle\langle (A^b) \rangle\rangle) \in q$, we have either $\psi \in q$ or $\psi \land \langle\langle (A^b) \rangle\rangle \in q$. Let us consider the following two sub-cases:

  - if $\psi \in q$, let us construct a simple tree $(T, V, C)$ where $T = \{e\}$ and $V(e) = q \cap cl(\phi_0) \supset \psi$.

    Then, $(T, V, C)$ satisfies all conditions of Lemma 5; hence, $q \cap cl(\phi_0) \in Z$; therefore, similar to the above argument of the case $\langle\langle (A^b) \rangle\rangle \psi \not\in q$, $\eta \in Z$.

  - if $\psi \not\in q$ and $\psi \land \langle\langle (A^b) \rangle\rangle \in q$. Let $\Phi = q \cap cl(\phi_0)$. Obviously, $\Phi \subseteq \Psi^* \subseteq \Psi^*$; therefore, $|\Phi| < k$. Then, by Lemma 4, there exists a locally consistent tree $(T_0, V_0, C_0)$ of height one where $T_0 = \{e\} \cup \{1, \ldots, k\}^n$ and $V_0(\epsilon) = q \cap cl(\phi_0)$. For each $e \in \{1, \ldots, k\}^n$, let $\Phi_e$ be an arbitrary set from $\Gamma$ such that $\Phi_e \supset V_0(\epsilon)$. Let $(T_1, V_1, C_1)$ be a finite tree such that $T_1 = T_0$, $C_1 = C_0$, $V_1(\epsilon) = V_0(\epsilon)$ and $V_1(\epsilon) = \Phi_e$ for all $e \in \{1, \ldots, k\}^n$. Since $(T_0, V_0, C_0)$ is locally consistent, so is $(T_1, V_1, C_1)$.

For every child $e \in \{1, \ldots, k\}^n$ such that $\chi \in V_1(e)$, we have:

$$\chi \in V_1(e) \Rightarrow V_1(e) \in Z \quad \text{by Lemma 5}$$

$$\Rightarrow \exists \text{ a local consistent tree } (T_e, V_e, C_e)$$

which satisfies all conditions of Lemma 6.

Let us consider a finite tree $(T, V, C)$ where

$$T = \{e\} \cup \{e \in \{1, \ldots, k\}^n | \chi \not\in V_1(e)\}$$

$$\{ c \in \{1, \ldots, k\}^n, t \in T_e, \chi \not\in V_1(c) \}$$

for all $e \in T$:

$$V(t) = \begin{cases} 
V_1(\epsilon) & \text{if } t = \epsilon \\
V_1(c) & \text{if } t = c \text{ where } c \in \{1, \ldots, k\}^n, \chi \not\in V_1(c) \\
V_e(\epsilon') & \text{if } t = \epsilon' \text{ where } c \in \{1, \ldots, k\}^n, t \in T_e, \chi \not\in V_1(c) 
\end{cases}$$
for all $t \in T$, $i \in N$ and $j \in \{1, \ldots, k\}$:

$$C(t, i, j) = \begin{cases} 
C_1(\epsilon, i, j) & \text{if } t = \epsilon \\
C_1(ct', i, j) & \text{if } t = ct' \text{ where } c \in \{1, \ldots, k\}^n, t' \in T_c, \chi \in V_1(c)
\end{cases}$$

It is straightforward that $(T, V, C)$ is also locally consistent and all of its interior nodes have $k^n$ children.

Let us show that $(T, V, C)$ realizes $\langle(A(\bar{A} B)) \psi U \psi \rangle$ at $\epsilon$. Let $c \in \{1, \ldots, k\}^n$ such that $C(c, A) \leq \bar{A} B c$, i.e. $c_A$ is a joint move which costs at most $\bar{A} B c$, we have:

$$\langle(A(\bar{A} B)) \chi Z \in V_1(\epsilon) = V(\epsilon)$$

$$\Rightarrow \forall c' \in \{1, \ldots, k\}^n: c'_A = c_A \Rightarrow \chi \in V(c)$$

$$\Rightarrow \forall c' \in \{1, \ldots, k\}^n: c'_A = c_A \Rightarrow V(c') = V(\epsilon) \in Z$$

$$\Rightarrow \forall c' \in \{1, \ldots, k\}^n: c'_A = c_A \Rightarrow (\langle(A(\bar{A} B)) \psi U \psi \rangle$$

It is straightforward that $F_A$ realizes $\langle(A(\bar{A} B)) \psi U \psi \rangle$ at the root of $(T_c, V_c, C_c)$.

Let us consider a $\bar{A} B$-strategy $F_A$ where

$$F_A(\lambda) = \begin{cases} 
c_A & \text{if } \lambda = \epsilon \\
F_A, c'(c't) & \text{if } c' \in \{1, \ldots, k\}^n, t \in T_c, \chi \in V_1(c')
\end{cases}$$

It is straightforward that $F_A$ realizes $\langle(A(\bar{A} B)) \psi U \psi \rangle$ from the root $\epsilon$ of $(T, V, C)$. Hence, as $T(\epsilon) = q \cap cl(\varphi_0)$, we have:

$$q \cap cl(\varphi_0) \in Z \Rightarrow \chi \in q \cap cl(\varphi_0) \quad \text{by Lemma 5}$$

$$\Rightarrow \chi \in q \quad \text{as } q \text{ is maximal}$$

**Induction step:** Assume that $b > \bar{A} B$. Similar to the base case, as $\psi \lor (\varphi \land (\langle A B \rangle) \psi U \psi \lor ((\langle(A(\bar{A} B)) \chi Z \rangle)) \in q$, let us consider the following three sub-cases:

- if $\varphi \in q$, the proof is the repetition of that for the base case.
- if $\varphi$ and $\langle(A(\bar{A} B)) \chi Z \in q$, the proof is the repetition of that for the base case.
- if $\varphi$ and $\langle(A(\bar{A} B)) \psi U \psi \in q$ for some $(b_1, b_2) \in split(b)$, let $\Phi = q \cap cl(\varphi_0)$. Obviously, $\Phi \subseteq \Psi \circ \Phi$, therefore, $\Phi \equiv \Phi$. By Lemma 4, there exists a locally consistent tree $(T_0, V_0, C_0)$ of height one where $T_0 = \{\epsilon\} \cup \{1, \ldots, k\}^n$ and $V_0(\epsilon) = \Phi$. For each $c \in \{1, \ldots, k\}^n$, let $\Phi_c$ be an arbitrary set from $\Gamma$ such that $\Phi_c \supseteq V_0(c)$. Let $(T_1, V_1, C_1)$ be a finite tree such that $T_1 = T_0$, $C_1 = C_0$, $V_1(\epsilon) = V_0(\epsilon)$ and $V_1(c) = \Phi_c$ for all $c \in \{1, \ldots, k\}^n$. Since $(T_0, V_0, C_0)$ is locally consistent, so is $(T_1, V_1, C_1)$.
For each $c \in \{1, \ldots, k\}^n$ such that $\langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi \in V(c)$, as $b_2 < b$, we have:

$$\langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi \in V(c) \Rightarrow \exists \text{ a locally consistent tree } (T_c, V_c, C_c)
$$

which realizes $\langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi$

by induction hypothesis

Let us consider a finite tree $(T, V, C)$ where

$$T = \{c\} \cup \{c \in [1, \ldots, k]^n : \langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi \notin V_1(c)\}
\cup \{ct \mid c \in [1, \ldots, k]^n, t \in T_c, \langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi \in V_1(c)\}
$$

for all $t \in T$:

$$V(t) = \begin{cases} V_1(\epsilon) & \text{if } t = \epsilon \\ V_1(c) & \text{if } t = c \text{ where } c \in [1, \ldots, k]^n, \langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi \notin V_1(c) \\ V_c(ct') & \text{if } t = ct' \text{ where } c \in [1, \ldots, k]^n, t' \in T_c, \langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi \in V_1(c) \end{cases}$$

for all $t \in T$, $i \in N$ and $j \in \{1, \ldots, k\}$:

$$C(t, i, j) = \begin{cases} C_1(\epsilon, i, j) & \text{if } t = \epsilon \\ C_c(ct', i, j) & \text{if } t = ct' \text{ where } c \in [1, \ldots, k]^n, t' \in T_c, \langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi \in V_1(c) \end{cases}$$

It is straightforward that $(T, V, C)$ is also locally consistent and all of its interior nodes have $k^n$ children.

Let us show that $(T, V, C)$ realizes $\langle \langle A^{b_1} \rangle \rangle \varphi \mathcal{U} \psi$ at $\epsilon$. Let $c \in [1, \ldots, k]^n$ such that $C(c, A) \leq b_1$, i.e., $c_A$ is a joint move which costs at most $b_1$, we have:

$$\langle \langle A^{b_1} \rangle \rangle \circ \langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi \in V_1(\epsilon) = V(\epsilon)
\Rightarrow \forall c' \in [1, \ldots, k]^n : c'_A = c_A \Rightarrow \langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi \in V(c')
\Rightarrow \forall c' \in [1, \ldots, k]^n : c'_A = c_A \Rightarrow V(c') = V(c) = Z
\Rightarrow \forall c' \in [1, \ldots, k]^n : c'_A = c_A \Rightarrow \langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi
\text{ is realized at the root of } (T_c, V_c, C_c)
\Rightarrow \forall c' \in [1, \ldots, k]^n : \exists b_2\text{-strategy } F_{A,c'}
\text{ realizes } \langle \langle A^{b_2} \rangle \rangle \varphi \mathcal{U} \psi \text{ at the root of } (T_c, V_c, C_c)
$$

Let us consider a $b$-strategy $F_A$ where

$$F_A(\lambda) = \begin{cases} c_A & \text{if } \lambda = \epsilon \\ F_{A,c'}(c't) & \text{if } c' \in [1, \ldots, k]^n, t \in T_{c'}, \exists Z \in V_1(c') \end{cases}$$
It is straightforward that $F_A$ realizes $\langle\langle A^b \rangle\rangle \varphi U \psi$ from the root $e$ of $(T, V, C)$. Hence, as $T(e) = q \cap cl(\varphi_0)$, we have:

\[
q \cap cl(\varphi_0) \in Z \Rightarrow \chi_Z \in q \cap cl(\varphi_0) \quad \text{by Lemma 5}
\]

\[
\Rightarrow \chi_Z \in q
\]

\[
\Rightarrow \varphi \in q \quad \text{as } q \text{ is maximal}
\]

\[\blacksquare\]

Similarly, we have the following result:

**Lemma 7**

Given $\neg\langle\langle A^b \rangle\rangle \Box \varphi$ and $x \in \Gamma$, there is a finite tree $(T, V, C)$ over $\Gamma$ such that:

- every interior node of $(T, V, C)$ has $k^n$ children where $k = |\Psi_0| + 1$
- $(T, V, C)$ is locally consistent
- $V(e) = x$
- if $\neg\langle\langle A^b \rangle\rangle \Box \varphi \in x$ then $(T, V, C)$ realizes $\neg\langle\langle A^b \rangle\rangle \Box \varphi$ from $e$

**Proof.** The proof is similar to that of the previous lemma. Let $Z \subseteq \Gamma$ such that, for any $x \in Z$, there is a finite tree obeying all conditions of Lemma 7. We shall prove the lemma by showing that $Z = \Gamma$.

In the following, assume that $x \in \Gamma$.

- If $\neg\langle\langle A^b \rangle\rangle \Box \varphi \notin x$, construct a simple tree $(T, V, C)$ where $T = \{e\}$ and $V(e) = x$. Since $(T, V, C)$ satisfies all conditions of the lemma, we have:

\[
q \cap cl(\varphi_0) \in Z \Rightarrow \chi_Z \in q \cap cl(\varphi_0) \quad \text{by Lemma 5}
\]

\[
\Rightarrow \chi_Z \in q
\]

\[
\Rightarrow \varphi \in q \quad \text{as } q \text{ is maximal}
\]

- If $\neg\langle\langle A^b \rangle\rangle \Box \varphi \in x$, we first show that $\eta = (\neg \varphi \lor (\neg\langle\langle A^b \rangle\rangle \Box \varphi \land \neg\langle\langle A^\infty \prec b \rangle\rangle \Box \neg \chi_Z)) \rightarrow \chi_Z$ is a theorem. Then, we have:

\[
\neg\langle\langle A^b \rangle\rangle \Box \varphi \in x \Rightarrow \eta \in x \text{ as } x \text{ is maximal}
\]

\[
\Rightarrow x \in Z \quad \text{by Lemma 5}
\]

However, it remains to prove that $\eta$ is a theorem. Again, this is done by showing that $\eta$ belongs to any maximal consistent set $q$ of RB-ATL in three cases:

- If $\neg\langle\langle A^b \rangle\rangle \Box \varphi \notin q$, let us construct a simple tree $(T, V, C)$ where $T = \{e\}$ and $V(e) = q \cap cl(\varphi_0)$. Since $(T, V, C)$ satisfies all conditions of the lemma, we have:

\[
q \cap cl(\varphi_0) \in Z \Rightarrow \chi_Z \in q \cap cl(\varphi_0) \quad \text{by Lemma 5}
\]

\[
\Rightarrow \chi_Z \in q
\]

\[
\Rightarrow \eta \in q \quad \text{as } q \text{ is maximal}
\]

- If $\neg \varphi \lor (\neg\langle\langle A^b \rangle\rangle \Box \varphi \land \neg\langle\langle A^\infty \prec b \rangle\rangle \Box \neg \chi_Z) \notin q$, then it is straightforward that $\eta \in q$ since $q$ is maximal.
• If \( \neg(\langle A^b \rangle \square \varphi) \) and \( \neg \varphi \lor (\neg(\langle A^{\text{inf}} \rangle) \square \varphi \land \neg(\langle A^{\text{inf}} \rangle) \square \neg \chi_Z) \in q \), we prove that \( \eta \in q \) by induction on \( b \).

**Base case:** Assume that \( b = 0^{\infty} \). As \( \neg \varphi \lor (\neg(\langle A^{\text{inf}} \rangle) \square \varphi \land \neg(\langle A^{\text{inf}} \rangle) \square \neg \chi_Z) \in q \), we have either \( \neg \varphi \in q \) or \( \neg(\langle A^{\text{inf}} \rangle) \square \neg \chi_Z \in q \). Let us consider the following two sub-cases:

- If \( \neg \varphi \in q \), let us construct a simple tree \((T, V, C)\) where \( T = \{ \epsilon \} \) and \( V(\epsilon) = q \land c(\varphi_0) \lor \neg \varphi \).
  Then, \((T, V, C)\) satisfies all conditions of Lemma 5; hence, \( q \land c(\varphi_0) \in Z \); therefore, similar to the above argument of the case \( \neg(\langle A^b \rangle) \square \varphi \notin q \), \( \eta \in Z \).

- If \( \neg(\langle A^{\text{inf}} \rangle) \square \neg \chi_Z \), let \( \Phi = q \land c(\varphi_0) \). Obviously, \( \Phi \subseteq \Psi \); therefore, \( |\Phi| < k \).
  Then, by Lemma 4, there exists a locally consistent tree \((T_0, V_0, C_0)\) of height one where \( T_0 = \{ \epsilon \} \cup \{ 1, \ldots, k \} \) and \( V_0(\epsilon) = q \land c(\varphi_0) \).
  For each \( c \in \{ 1, \ldots, k \} \), let \( \Phi_c \) be an arbitrary set from \( \Gamma \) such that \( \Phi_c \supseteq V_0(c) \).
  Let \((T_1, V_1, C_1)\) be a finite tree such that \( T_1 = T_0, C_1 = C_0, V_1(\epsilon) = V_0(\epsilon) \) and \( V_1(c) = \Phi_c \) for all \( c \in \{ 1, \ldots, k \} \).
  Since \((T_0, V_0, C_0)\) is locally consistent, so is \((T_1, V_1, C_1)\).
  Then, for every \( c \in \{ 1, \ldots, k \} \) such that \( C(c, A) \leq 0 \leq b \), i.e., \( c_A \) is a joint move which costs at most \( 0 \leq b \), we have:

\[
C(c, A) \leq 0 \leq b \Rightarrow \exists c' \{ 1, \ldots, k \} : c'_A = c_A \land \\
\neg \chi_Z \lor \chi_Z \in V_1(c') \\
\Rightarrow V_1(c') \in Z \text{ by Lemma 5} \\
\Rightarrow \exists \text{ a locally consistent tree } (T_{c'}, V_{c'}, C_{c'}) \\
\text{ which satisfies all conditions of Lemma 7}
\]

Let us consider a finite tree \((T, V, C)\) where

\[
T = \{ \epsilon \} \cup \{ c \in \{ 1, \ldots, k \} \mid \chi_Z \notin V_1(c) \} \\
\{ ct \mid c \in \{ 1, \ldots, k \}, t \in T_{c}, \chi_Z \in V_1(c) \}
\]

for all \( t \in T \):

\[
V(t) = \begin{cases} 
V_1(\epsilon) \text{ if } t = \epsilon \\
V_1(\epsilon) \text{ if } t = c \text{ where } c \in \{ 1, \ldots, k \}, \chi_Z \notin V_1(c) \\
V_{c}(ct') \text{ if } t = ct' \text{ where } c \in \{ 1, \ldots, k \}, t' \in T_{c}, \chi_Z \in V_1(c)
\end{cases}
\]

for all \( t \in T, i \in N \) and \( j \in \{ 1, \ldots, k \} \):

\[
C(t, i, j) = \begin{cases} 
C_{1}(\epsilon, i, j) \text{ if } t = \epsilon \\
C_{c}(ct', i, j) \text{ if } t = ct' \text{ where } c \in \{ 1, \ldots, k \}, t' \in T_{c}, \chi_Z \in V_1(c)
\end{cases}
\]

It is straightforward that \((T, V, C)\) is also locally consistent and all of its interior nodes have \( k^n \) children.
Let us show that \((T, V, C)\) realizes \(\langle A^{\infty_b} \rangle \boxslash \varphi\) at \(\epsilon\). Let \(c \in [1, \ldots, k]^n\) such that \(C(c, A) \leq 0 \overset{\infty}{\rightarrow} b\), i.e. a joint move that costs at most \(0 \overset{\infty}{\leftarrow} b\), we have:

\[
\neg\langle A^{\infty_b} \rangle \circ \neg C \in V_1(\epsilon) \Rightarrow V(\epsilon)
\]

\[
\Rightarrow \exists c' \in [1, \ldots, k]^n: c' = c_A \wedge C \in V(c')
\]

\[
\Rightarrow V(c') = V_c(\epsilon) \in Z
\]

\[
\Rightarrow \neg\langle A^{\infty_b} \rangle \boxslash \varphi
\]

is realized at the root of \((T_c, V_c, C_c)\)

\[
\Rightarrow \neg\langle A^{\infty_b} \rangle \boxslash \varphi
\]

is realized at the root of \((T, V, C)\)

Hence, as \(T(\epsilon) = q \cap cl(\varphi_0)\), we have:

\[
q \cap cl(\varphi_0) \in Z \Rightarrow C \in q \cap cl(\varphi_0) \quad \text{by Lemma 5}
\]

\[
\Rightarrow C \in q
\]

\[
\Rightarrow q \in q \quad \text{as } q \text{ is maximal}
\]

**Induction step:** Assume that \(b > 0 \overset{\infty}{\rightarrow} \infty\). Similar to the base case, as \(\neg \varphi \lor (\neg\langle A^b \rangle \circ \Box \varphi \land \neg\langle A^{\infty_b} \rangle \circ \neg C) \in q\), let us consider the following two sub-cases:

- if \(\neg \varphi \in q\), the proof is the repetition of that for the base case.
- if \(\neg\langle A^b \rangle \circ \Box \varphi \lor \neg\langle A^{\infty_b} \rangle \circ \neg C \in q\), then \(\neg\langle A^{\infty_b} \rangle \circ \neg C \in q\) and \(\neg\langle A^b \rangle \circ \Box \varphi \lor \neg\langle A^{\infty_b} \rangle \circ \neg C \in q\) for all \((b_1, b_2) \in \text{split}(b)\). Let \(\Phi = q \cap cl(\varphi_0)\). Obviously, \(\Phi \subseteq \Psi\); therefore, \(\Phi \subseteq \kappa\). By Lemma 4, there exists a local consistent tree \((T_0, V_0, C_0)\) of height one where \(T_0 = \{\epsilon\} \cup \{1, \ldots, k\}^n\) and \(V_0(\epsilon) = \Phi\).

For each \(c \in [1, \ldots, k]^n\), let \(\Phi_c\) be an arbitrary set from \(\Gamma\) such that \(\Phi_c \supseteq V_0(\epsilon)\). Let \((T_1, V_1, C_1)\) be a finite tree such that \(T_1 = T_0, C_1 = C_0, V_1(\epsilon) = V_0(\epsilon)\) and \(V_1(\epsilon) = \Phi_c\) for all \(c \in [1, \ldots, k]^n\). Since \((T_0, V_0, C_0)\) is locally consistent, so is \((T_1, V_1, C_1)\). Then, for every \(c \in [1, \ldots, k]^n\) such that \(C(c, A) \leq b\), i.e., \(c_A\) is a joint move which costs at most \(b\), we have:

if \(C(c, A) \leq 0 \overset{\infty}{\rightarrow} b\):

\[
C(c, A) \leq \tilde{b} \Rightarrow \exists c' [1, \ldots, k]^n: c' = c_A \land \neg \neg C \in V_1(c')
\]

\[
\Rightarrow V_1(c') \in Z \quad \text{by Lemma 5}
\]

\[
\Rightarrow \exists \text{ a locally consistent tree } (T_c, V_c, C_c)
\]

which satisfies all conditions of Lemma 7

if \(0 \overset{\infty}{\rightarrow} b < C(c, A) \leq b_1\) for some \((b_1, b_2) \in \text{split}(b)\):

\[
C(c, A) \leq \tilde{b} \Rightarrow \exists c' [1, \ldots, k]^n: c' = c_A \land \neg \langle A^{b_1} \rangle \Box \varphi \in V_1(c')
\]

\[
\Rightarrow V_1(c') \in Z \quad \text{induction hypothesis}
\]

\[
\Rightarrow \exists \text{ a locally consistent tree } (T_c, V_c, C_c)
\]

which satisfies all conditions of Lemma 7
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Let us consider a finite tree \((T, V, C)\) where

\[
T = \{\varepsilon\} \cup \{c \in \{1, \ldots, k\}^n \mid \chi_Z \not\in V_1(c)\} \cup \\
\{ct \mid c \in \{1, \ldots, k\}^n, t \in T_c, \chi_Z \in V_1(c) \text{ or} \\
\exists (b_1, b_2) \in \text{split}(b) : \neg\langle\langle A^{b_2} \rangle\rangle \square \varphi \in V_1(c)\}
\]

for all \(t \in T\):

\[
V(t) = \begin{cases} 
V_1(\epsilon) & \text{if } t = \epsilon \\
V_1(c) & \text{if } t = c \text{ where} \\
V_c(\epsilon) & \text{if } t = c \epsilon' \text{ where} \\
& \quad c \in \{1, \ldots, k\}^n, \chi_Z \not\in V_1(c) \\
& \quad c \in \{1, \ldots, k\}^n, t' \in T_c, \chi_Z \in V_1(c) \text{ or} \\
& \quad \exists (b_1, b_2) \in \text{split}(b) : \neg\langle\langle A^{b_2} \rangle\rangle \square \varphi \in V_1(c) \end{cases}
\]

and for all \(t, i \in N\) and \(j \in \{1, \ldots, k\}\):

\[
C(t, i, j) = \begin{cases} 
C_1(\epsilon, i, j) & \text{if } t = \epsilon \\
C_c(\epsilon', i, j) & \text{if } t = c \epsilon' \text{ where} \\
& \quad c \in \{1, \ldots, k\}^n, t' \in T_c, \chi_Z \in V_1(c) \text{ or} \\
& \quad \exists (b_1, b_2) \in \text{split}(b) : \neg\langle\langle A^{b_2} \rangle\rangle \square \varphi \in V_1(c) \end{cases}
\]

It is straightforward that \((T, V, C)\) is also locally consistent and all of its interior nodes have \(k^n\) children.

Let us show that \((T, V, C)\) realizes \(\langle\langle A^1 \rangle\rangle \square \varphi\) at \(\varepsilon\). Let \(c \in \{1, \ldots, k\}^n\) such that \(C(c, A) \leq \tilde{b}\), i.e. a joint move which costs at most \(b\), we have:

if \(C(c, A) \leq \tilde{b}\):

\[
\neg\langle\langle A^{\tilde{b}} \rangle\rangle \circ \neg\chi_Z \in V_1(\epsilon) = V(\epsilon) \\
\Rightarrow \exists c' \in \{1, \ldots, k\}^n : \chi_{c'} = \chi_A \Rightarrow \chi_Z \in V(c') \\
\Rightarrow V(c') = V_c(\epsilon) \in Z \\
\Rightarrow \neg\langle\langle A^{\tilde{b}} \rangle\rangle \square \varphi \text{ is realized at the root of } (T_{c'}, V_{c'}, C_{c'}) \\
\Rightarrow \neg\langle\langle A^{\tilde{b}} \rangle\rangle \square \varphi \text{ is realized at the root of } (T, V, C)
\]

if \(\tilde{b} < C(c, A) \leq b_1\) for some \((b_1, b_2) \in \text{split}(b)\):

\[
\neg\langle\langle A^{b_1} \rangle\rangle \circ \langle\langle A^{b_2} \rangle\rangle \square \varphi \in V_1(\epsilon) = V(\epsilon) \\
\Rightarrow \exists c' \in \{1, \ldots, k\}^n : \chi_{c'} = \chi_A \Rightarrow \neg\langle\langle A^{b_2} \rangle\rangle \square \varphi \in V(c') \\
\Rightarrow V(c') = V_{c'}(\epsilon) \in Z \\
\Rightarrow \neg\langle\langle A^{\tilde{b}} \rangle\rangle \square \varphi \text{ is realized at the root of } (T_{c'}, V_{c'}, C_{c'}) \\
\Rightarrow \neg\langle\langle A^{\tilde{b}} \rangle\rangle \square \varphi \text{ is realized at the root of } (T, V, C)
\]
Hence, as $T(\varepsilon) = q \cap cl(\varphi_0)$, we have:

$$q \cap cl(\varphi_0) \in \mathcal{Z} \Rightarrow \chi_{\mathcal{Z}} \in q \cap cl(\varphi_0) \quad \text{by Lemma 5}$$

$$\Rightarrow \chi_{\mathcal{Z}} \in q$$

$$\Rightarrow \eta \in q \quad \text{as } q \text{ is maximal}$$

Now we have almost all the ingredients for constructing the model for $\varphi_0$. For each consistent set $x$ in $\Gamma$ and an eventuality $\varphi$ of $cl(\varphi_0)$, we have a finite tree $(T_x, \varphi, V_x, \pi, C_x, \varphi)$ with the root having label $x$ which realizes $\varphi$. Let the eventualities in $cl(\varphi_0)$ be listed as $\varphi_1', \ldots, \varphi_m'$. Next, we define the final tree.

**Definition 8**

The final tree $(T_{\varphi_0}, V_{\varphi_0}, C_{\varphi_0})$ is constructed inductively as follows.

- Initially, select an arbitrary $x \in \Gamma$ such that $\varphi_0 \in x$. As that formula is consistent, such a set exists.
- Let $(T_x, \varphi_0', V_x, \varphi_0', C_x, \varphi_0')$ be the initial tree.
- Given the tree constructed so far and the last used eventuality $\varphi_j'$. We replace every leaf labelled by $y \in \Gamma$ of the currently constructed tree with the tree $(T_y, \varphi_j', V_y, \varphi_j', C_y, \varphi_j')$ where $j = i \mod m + 1$.

Let $S_{\varphi_0}$ be the model which is based on $(T_{\varphi_0}, V_{\varphi_0}, C_{\varphi_0})$. (It is easy to define the assignment $\pi$ using $V_x$.)

**Lemma 8**

If $\langle \langle A^b \rangle \rangle \varphi \mathcal{U} \psi$ (or $\neg \langle \langle A^b \rangle \rangle \square \varphi$) is in the label of some $t$ of $(T_{\varphi_0}, V_{\varphi_0}, C_{\varphi_0})$, $\langle \langle A^b \rangle \rangle \varphi \mathcal{U} \psi$ (or $\neg \langle \langle A^b \rangle \rangle \square \varphi$) is realized from $t$.

**Proof.** Let us consider the first case when $\varphi_i' = \langle \langle A^b \rangle \rangle \varphi \mathcal{U} \psi \in V(t)$ where $t$ is a node of $(T_{\varphi_0}, V_{\varphi_0}, C_{\varphi_0})$. The proof for the case of $\varphi_i' = \neg \langle \langle A^b \rangle \rangle \square \varphi$ is also done similarly.

- If $t$ happens to be the root of the sub-tree $(T_{\varphi_i'}, V_{\varphi_i'}, C_{\varphi_i'})$, then the proof is done as $\varphi_i'$ is realized within this sub-tree at $t$, hence also in the final tree.
- Otherwise, we define inductively on $b$ a $b$-strategy as follows:
  
  **Base case:** Assume that $b = 0 \stackrel{\infty}{\longrightarrow} b$, since $\langle \langle A^0 \stackrel{\infty}{\longrightarrow} b \rangle \rangle \varphi \mathcal{U} \psi \in V(t)$, as $V(t)$ is a maximally consistent set, we have that $\varphi \lor (\varphi \land \langle \langle A^0 \stackrel{\infty}{\longrightarrow} b \rangle \rangle \varphi \mathcal{U} \psi) \in V(t)$:
  
  - If $\psi \in V(t)$, the proof is done as $\langle \langle A^0 \stackrel{\infty}{\longrightarrow} b \rangle \rangle \varphi \mathcal{U} \psi$ is immediately realized at $t$.
  - Otherwise, we have $\varphi \land \langle \langle A^0 \stackrel{\infty}{\longrightarrow} b \rangle \rangle \varphi \mathcal{U} \psi \in V(t)$. Then $\varphi \in V(t)$ and by Lemma 4, there exists $c \in \{1, \ldots, k\}^n$ such that $C(ic, A) \leq 0 \stackrel{\infty}{\longrightarrow} b$ and for all $c \in \{1, \ldots, k\}^n$ with $c_A = c_{A'}$, we have $\langle \langle A^0 \stackrel{\infty}{\longrightarrow} b \rangle \rangle \varphi \mathcal{U} \psi \in V(tc')$. Let $F_A(t) = c_A$. Then, we can continue with the same argument to define the strategy $F_A$ until a node $t'$ labelled by $y$ in $(T_{\psi_y}, V_{\psi_y}, C_{\psi_y})$ is reached where $t'$ is the root of some sub-tree $(T_{\psi}', V_{\psi}', C_{\psi'})$. Such a node must exist because, according to the construction of the $(T_{\psi_y}, V_{\psi_y}, C_{\psi_y})$, eventual formulas in $cl(\varphi_0)$ are cycled through. As $(T_{\psi}', V_{\psi}', C_{\psi'})$ realizes $\varphi_i'$, we can extend $F_A$ to a $b$-strategy to realize $\varphi_i'$.

**Induction Step:** Assume that $b > 0 \stackrel{\infty}{\longrightarrow} b$, since $\langle \langle A^b \rangle \rangle \varphi \mathcal{U} \psi \in V(t)$, and $V(t)$ is a maximally consistent set, we have that $\varphi \lor (\varphi \land \langle \langle A^b \rangle \rangle \varphi \mathcal{U} \psi \lor \langle \langle A^0 \stackrel{\infty}{\longrightarrow} b \rangle \rangle \varphi \mathcal{U} \psi) \in V(t)$:

- If $\psi \in V(t)$, the proof is done as $\langle \langle A^b \rangle \rangle \varphi \mathcal{U} \psi$ is immediately realized at $t$. 

[16:41 5/6/2015 exv034.tex] LogCom: Journal of Logic and Computation
Then, we define \( \lambda \) all

\[ \langle \langle \text{is realised from } \rangle \rangle \]

Lemma 9

\[ \langle \langle \text{Induction step: } \rangle \rangle \]

In the following, let \( \phi \), \( \lambda = \langle \langle \rangle \rangle \), and by Lemma 4 and there exists \( \phi \) and \( \langle \langle \phi \rangle \rangle \) from \( t \). Hence, we just need to define \( F_A(t) = \lambda \). As \( \phi \) and \( \langle \langle \phi \rangle \rangle \) from \( t \).

Otherwise, we have \( \phi \) and \( \langle \langle \phi \rangle \rangle \) from \( \langle \langle \phi \rangle \rangle \) is realised from \( t \). Let us repeat the argument in the base case where \( \phi \) is realised from \( t \) and by Lemma 4 and we have that there exists \( c \in \{ 1, \ldots, k \} \) such that \( R(t, c) \leq b \) and for all \( c \in \{ 1, \ldots, k \} \) such that \( c^\prime = c_A \), we have \( \langle \langle \phi \rangle \rangle \) from \( t \). Let \( F_A(t) = c_A \). Then, we can continue with the same argument to define the strategy \( F_A \) until a node \( t^\prime \) labelled by \( y \) in \((T_{v_0}, V_{v_0}, C_{v_0})\) is reached where \( t^\prime \) is the root of some sub-tree \((T_{v_1}, V_{v_1}, C_{v_1})\). Such a node must exist because, according to the construction of the \((T_{v_0}, V_{v_0}, C_{v_0})\), eventualities in \( c(t_{v_0}) \) are cycled through. As \((T_{v_1}, V_{v_1}, C_{v_1})\) realises \( \phi \), we can extend \( F_A \) to a \( b \)-strategy to realize \( \phi \).

\[ \text{Lemma 9} \]

If \( \neg \langle \langle \phi \rangle \rangle \text{ is in the label of some } t \) of \((T_{v_0}, V_{v_0}, C_{v_0})\), \( \neg \langle \langle \phi \rangle \rangle \text{ is realised from } t \).

**Proof.** Let us consider the case \( \langle \langle \phi \rangle \rangle \text{ is in the label of some } t \) of \((T_{v_0}, V_{v_0}, C_{v_0})\). The proof for \( \neg \langle \langle \phi \rangle \rangle \text{ is realised from } t \) is similar.

In the following, let \( t^\prime = \{ 1, \ldots, k \} \). We shall define a \( b \)-strategy \( F_A \) which realizes \( \langle \langle \phi \rangle \rangle \text{ is in the label of some } t \) of \((T_{v_0}, V_{v_0}, C_{v_0})\). Such a node must exist because, according to the construction of the \((T_{v_0}, V_{v_0}, C_{v_0})\), eventualities in \( c(t_{v_0}) \) are cycled through. As \((T_{v_1}, V_{v_1}, C_{v_1})\) realises \( \phi \), we can extend \( F_A \) to a \( b \)-strategy to realize \( \phi \).

**Base case:** Assume that \( i = 1 \). We have

\[ \langle \langle \phi \rangle \rangle \in V(t) \Rightarrow \phi \in V(t) \land \langle \langle \phi \rangle \rangle \in V(t) \]

for some \((b_1, b_2) \in \text{split}(b) \cup \{(0 \infty b, b)\})

\[ \Rightarrow \exists c \in t^\prime: C(t, c) \leq b_1 \]

\[ \Rightarrow \forall c^\prime \in t^\prime: c^\prime = c_A \Rightarrow \langle \langle \phi \rangle \rangle \text{ in } V(t^\prime) \]

Then, we define \( F_A(t) = c_A \) and \( \lambda^2 = \{ t^\prime \} \). Obviously, we have \( \text{cost}(t^\prime, F_A) = C(t^\prime, A) \leq b_1 \) and \( \langle \langle \phi \rangle \rangle \text{ in } V(t^\prime) \) for all \( t^\prime \in \lambda^2 \).

**Induction step:** Assume that \( i > 1 \) and we have defined \( F_A \) for inputs of length \( i \) and \( \lambda = \langle \langle \rangle \rangle \) such that for all \( \lambda \in \lambda^i \), \( \text{cost}(\lambda, F_A) \leq b_1 \) and \( \langle \langle \phi \rangle \rangle \text{ in } V(\lambda[i+1]) \) for some \((b_1, b_2) \in \text{split}(b) \cup \{(0 \infty b, b)\}) \). Let us define \( F_A \) for inputs of length \( i+1 \) and \( \lambda^i \). We have, for all \( \lambda \in \lambda^i + 1 \)

\[ \langle \langle \phi \rangle \rangle \text{ in } V(\lambda[i+1]) \Rightarrow \phi \in V(\lambda[i+1]) \land \langle \langle \phi \rangle \rangle \text{ in } V(\lambda[i+1]) \]

for some \((b_3, b_4) \in \text{split}(b) \cup \{(0 \infty b, b)\})

\[ \Rightarrow \exists c_3 \in t^\prime: C(c, A) \leq b_3 \]

\[ \Rightarrow \forall c^\prime \in t^\prime: c^\prime = c_A \Rightarrow \langle \langle \phi \rangle \rangle \text{ in } V(\lambda(c^\prime)) \]

Then, we define \( F_A(\lambda) = c_A \) for all \( \lambda \in \lambda^i + 1 \) and \( \lambda^i + 2 = \{ \lambda^i + 1, c^\prime \} \). Obviously, we have \( (b_1 + b_2 + b_3) \in \text{split}(b) \cup \{(0 \infty b, b)\} \), \( \text{cost}(\lambda, F_A) = \text{cost}(\lambda, F_A) \) and \( \langle \langle \phi \rangle \rangle \in V(\lambda(c^\prime)) \) for all \( \lambda c^\prime \in \lambda^2 \).
Let $\Lambda = \bigcup_{i \geq 0} \Lambda^i$, we have
\[
\forall \lambda \in \Lambda \Rightarrow \varphi \in V(\lambda[i]) \land \exists i \geq 0 : \lambda \in \Lambda^i \\
\Rightarrow \varphi \in V(\lambda[i]) \land \\
\exists (b_1, b_2) \in \text{split}(b) \cup \{(0, \infty b, b)\} \text{cost}(\lambda, F_A) \leq b_1 \\
\Rightarrow \varphi \in V(\lambda[i]) \land \text{cost}(\lambda, F_A) \leq b
\]

Given the constructed strategy $F_A$, we have that $\text{out}(t, F_A) = \Lambda$. As $\text{cost}(\lambda, F_A) \leq b$ for all $\lambda \in \Lambda$, $F_A$ is a $b$-strategy. Furthermore, as $\varphi \in V(\lambda[i])$ for all $\lambda \in \Lambda$, it is straightforward that $F_A$ realizes $(A^b)\Box \psi$.

Finally, we show the following truth lemma.

**Lemma 10**
For every node $t$ of $(T_{\psi_0}, V_{\psi_0}, C_{\psi_0})$ and every formula $\varphi \in c(\psi_0)$, if $\varphi \in V_{\psi_0}(t)$ then $S_{\psi_0}, t \models \varphi$.

**Proof.** The proof is done by induction on the structure of $\varphi$.
- For the cases of propositions, negative proposition and disjunction, the proofs are trivial.
- Assume $\varphi = \langle \langle A^b \rangle \rangle \bigcirc \psi$, Lemma 4 ensures that there is a move $c_A$ for some $c \in \{1, \ldots, k\}^n$ where $C(c, A) \leq b$ such that for all $c' \in \{1, \ldots, k\}^n$, we have $\psi \in V(tc')$. Then by the induction hypothesis, we have that $S_{\psi_0}, tc' \models \psi$. Then, $S_{\psi_0}, t \models \langle \langle A^b \rangle \rangle \bigcirc \psi$.
- Assume $\varphi = \neg \langle \langle A^b \rangle \rangle \bigcirc \psi$, Lemma 4 ensures that for all $c \in \{1, \ldots, k\}^n$ such that $C(c, A) \leq b$, there exists $c' \in \{1, \ldots, k\}^n$ such that $c'_A = c_A$ and $\neg \psi \in V(tc')$. Then by the induction hypothesis, we have that $S_{\psi_0}, tc' \models \neg \psi$. Then, $S_{\psi_0}, t \models \neg \langle \langle A^b \rangle \rangle \bigcirc \psi$.
- For the cases of $\langle \langle A^b \rangle \rangle \psi_1 U \psi_2$, $\neg \langle \langle A^b \rangle \rangle \Box \psi$, $\langle \langle A^b \rangle \rangle \psi_1 U \psi_2$ and $\langle \langle A^b \rangle \rangle \Box \psi$, the proofs are trivial due to Lemmas 8 and 9.

Finally, we have the following theorem.

**Theorem 1**
The axiom system for RB-ATL is sound and complete.

### 4 Model-checking RB-ATL

In this section, we describe a model-checking algorithm for RB-ATL which runs in time polynomial in the size of the formula (if resource bounds are encoded in unary) and the structure, and is exponential in the number of resources. The algorithm is similar to the model-checking algorithm for ATL given in [6]. The main differences from the algorithm for ATL are that we need to take the costs of strategies into account, and, instead of working with a straightforward set of subformulas $\text{Sub}(\varphi)$ of a given formula $\varphi$, we work with an extended set of subformulas $\text{Sub}^+(\varphi)$. $\text{Sub}^+(\varphi)$ includes $\text{Sub}(\varphi)$, and in addition:

- if $\langle \langle A^d \rangle \rangle \Box \psi \in \text{Sub}(\varphi)$, then $\langle \langle A^{d'} \rangle \rangle \Box \psi \in \text{Sub}^+(\varphi)$ for all $d'$ such that $(d, d') \in \text{split}(b)$;
- if $\langle \langle A^d \rangle \rangle \psi_1 U \psi_2 \in \text{Sub}(\varphi)$, then $\langle \langle A^{d'} \rangle \rangle \psi_1 U \psi_2 \in \text{Sub}^+(\varphi)$ for all $d'$ such that $(d, d') \in \text{split}(b)$.

We assume that $\text{Sub}^+(\varphi)$ is ordered in the increasing order of complexity and of resource bounds (so e.g. if $b \leq b'$, then $\langle \langle A^b \rangle \rangle \Box \psi$ precedes $\langle \langle A^{b'} \rangle \rangle \Box \psi$).
Theorem 2
Given a structure $S=(n,r,Q,\Pi,\pi,d,e,\delta)$ and a formula $\varphi$, there is an algorithm which returns the set of states $[\varphi]_S$ satisfying $\varphi$: $[\varphi]_S = \{q \mid S.q \models \varphi\}$, which runs in time $O(|\varphi|^{2r+1} \times |S|)$, assuming resource bounds are encoded in unary.

Proof. Consider the following model-checking algorithm:

for every $\varphi'$ in $Sub^+(\varphi)$:
  case $\varphi'==p$: $[\varphi']_S = \pi(p)$
  case $\varphi'==\neg\psi$: $[\varphi']_S = Q \setminus [\psi]_S$
  case $\varphi'==\psi_1 \land \psi_2$: $[\varphi']_S = [\psi_1]_S \cap [\psi_2]_S$
  case $\varphi'==\langle\langle A^b \rangle\rangle \psi$: $[\varphi']_S = \text{Pre}(A,[\psi]_S,b)$
  case $\varphi'==\langle\langle A^b \rangle\rangle \Box \psi$ for $b$ where for all $i, b_i \in [0, \infty)$:
    $\rho := \text{[true]}_S$; $\tau := [\psi]_S$;
    while $\rho \not\subseteq \tau$ do $\rho := \tau$; $\tau := \text{Pre}(\rho, \rho, b) \cap [\psi]_S$ od;
    $[\varphi']_S := \rho$
  case $\varphi'==\langle\langle A^b \rangle\rangle \psi_1 U \psi_2$ for $b$ where for some $i, b_i \not\in [0, \infty)$:
    $\rho := \text{[false]}_S$; $\tau := [\psi_2]_S$;
    foreach $d' \in \{d' \mid (d,d') \in \text{split}(b)\}$ do
      $\tau := \text{Pre}(\rho, [\langle\langle A^d \rangle\rangle \psi_1 U \psi_2]_S, d) \cap [\psi]_S$
      while $\tau \not\subseteq \rho$ do $\rho := \rho \cup \tau$; $\tau := \text{Pre}(\rho, \rho, 0 \times b) \cap [\psi]_S$ od;
    end
    $[\varphi']_S := \rho$
  case $\varphi'==\langle\langle A^b \rangle\rangle \psi_1 U \psi_2$ where for some $i, b_i \not\in [0, \infty)$:
    $\rho := \text{[false]}_S$; $\tau := [\psi_2]_S$;
    foreach $d' \in \{d' \mid (d,d') \in \text{split}(b)\}$ do
      $\tau := \text{Pre}(\rho, [\langle\langle A^d \rangle\rangle \psi_1 U \psi_2]_S, d) \cap [\psi]_S$
      while $\tau \not\subseteq \rho$ do $\rho := \rho \cup \tau$; $\tau := \text{Pre}(\rho, \rho, 0 \times b) \cap [\psi]_S$ od;
    end
    $[\varphi']_S := \rho$

Pre is a function which given a coalition $A$, a set $\rho \subseteq Q$ and a bound $b$ returns a set of states $q$ in which $A$ has a move $\sigma_A$ with cost $\text{cost}(q, \sigma_A) \leq b$ such that $\text{out}(q, \sigma_A) \subseteq \rho$. Observe that $\text{Pre}(A, \rho, \infty)$ is just $\text{Pre}(A, \rho)$ from [6].

The cases for propositional variables, negation, conjunction and $\langle\langle A^b \rangle\rangle \psi$ are straightforward. The cases where the resource bound consists of 0 and $\infty$ are also similar to [6]. However, the cases for $\langle\langle A^b \rangle\rangle \Box \psi$ and $\langle\langle A^b \rangle\rangle \psi_1 U \psi_2$ where $b$ does not contain only 0 and $\infty$ have no counterpart in the ATL algorithm, and we explain these in some detail. First, note that the cases for $\langle\langle A^b \rangle\rangle \Box \psi$ and $\langle\langle A^b \rangle\rangle \psi_1 U \psi_2$ where $b$ consists of 0 and $\infty$ are the standard greatest and least fixed point computations, respectively, which consider only 0 cost moves for $i$ with $b_i=0$ or any moves for $i$ with $b_i=\infty$. In
A logic for describing Coalitional Resource Games and a model-checking procedure for the logic

All of these papers consider only the model-checking problem, and some of the logics allow both

requirements in these logics. The only work in this tradition that considered resources is [17], which

force the system into one of the

\[[\langle \bar{A}^{\infty}_{b} \rangle] \square \psi \] _S = Pre(A, [\langle \bar{A}^{\infty}_{b} \rangle] \square \psi ) _S, \bar{A}^{\infty}_{b} [\psi ]_{S}

and

\[[\langle \bar{A}^{\infty}_{b} \rangle] \psi _1 U \psi _2 ]_{S} = Pre(A, [\langle \bar{A}^{\infty}_{b} \rangle] \psi _1 U \psi _2 ]_{S}, \bar{A}^{\infty}_{b} [\psi ]_{S}

\[[\langle \bar{A}^{\infty}_{b} \rangle] \square \psi \] _S contains all states where A has a \bar{A}^{\infty}_{b}-cost strategy to maintain \psi forever. Note that A has a \bar{A}^{\infty}_{b}-cost strategy to maintain \psi forever if and only if it has a \bar{A}^{\infty}_{b}-cost strategy to force the system into one of the \[[\langle \bar{A}^{\infty}_{b} \rangle] \square \psi \] _S states, while maintaining \psi. In other words, to compute \[[\langle \bar{A}^{b} \rangle] \square \psi \] for \bar{A} containing \bar{A} \notin [0, \infty], we need to compute \[[\langle \bar{A}^{b} \rangle] \psi U \langle \bar{A}^{\infty}_{b} \rangle] \square \psi \]. This explains the similarity between the cases of \[[\langle \bar{A}^{b} \rangle] \square \psi \] and \[[\langle \bar{A}^{b} \rangle] \psi _1 U \psi _2 \] for the case of \bar{A} not consisting solely of 0 and \infty. In the case of \[[\langle \bar{A}^{b} \rangle] \square \psi \], in the first execution of the \textbf{foreach} \bar{d} \in \{\bar{d}' | (d, \bar{d}') \in split(\bar{b})\} loop, we have \bar{d}' = \bar{A}^{\infty}_{b} \bar{b} and \tau = Pre(A, [\langle \bar{A}^{\infty}_{b} \rangle] \square \psi ) _S, \bar{A}^{\infty}_{b} [\psi ]_{S}, which includes \[Pre(A, [\langle \bar{A}^{b} \rangle] \square \psi ) _S, \bar{A}^{\infty}_{b} [\psi ]_{S}, \] hence it also includes \[[\langle \bar{A}^{\infty}_{b} \rangle] \square \psi \] _S. In the nested \textbf{while} loop, \rho accumulates the results and \tau adds the \psi-states from where A has a \bar{A}^{\infty}_{b} strategy to enforce the outcome to be in \rho. In the outer loop, \bar{d} bounds are used in some order consistent with <, namely satisfying the condition that if \bar{b}_i < \bar{d}_j then \bar{b}_j is used before \bar{b}_i.

In the case for \[[\langle \bar{A}^{b} \rangle] \psi _1 U \psi _2 \] where \bar{A} does not consist only of 0 and \infty, after the first iteration of the \textbf{foreach} \bar{d} \in \{\bar{d}' | (d, \bar{d}') \in split(\bar{b})\} loop, \tau is \[[\langle \bar{A}^{\infty}_{b} \rangle] \psi _1 U \psi _2 ]_{S} which includes \[[\psi ]_{S}\]. The rest is very similar to the case for \[[\langle \bar{A}^{b} \rangle] \square \psi \] where \bar{A} does not consist solely of 0 and \infty.

Note that [split(\bar{b})] is \textbf{O}(b'). If \psi contains operators with bounds other than 0 and \infty, \textbf{[Sub^+}(\psi)]] is \textbf{O}(\psi \times |\psi'|'), assuming resource bounds are written in unary. In the \[[\langle \bar{A}^{b} \rangle] \square \psi \] and \[[\langle \bar{A}^{b} \rangle] \psi _1 U \psi _2 \] cases, the outer loop is executed \textbf{O}(\psi') times and the inner loop is executed in total at most |S| times. This gives us complexity \textbf{O}(\psi \times |\psi'| \times |\psi'| \times |S|), or \textbf{O}(\psi|^{2r+1} \times |S|). Note that the lower bound for model-checking complexity is given by the model-checking complexity of ATL, which is polynomial time in the size of the model and the formula.

5 Related work

Recent work on ATL and CL (e.g. [1, 6, 12, 13, 15, 16]) has allowed the expression of many interesting

properties of coalitions and strategies. However, there is no natural way of expressing resource

requirements in these logics. The only work in this tradition that considered resources is [17], which

introduced Coalitional Resource Games and studied complexity of decision problems for these games.

A logic for describing Coalitional Resource Games and a model-checking procedure for the logic

were proposed in [3]; however, the only modality that logic has is \[[\langle \bar{A}^{b} \rangle] \square \) (only one-step games

were considered).

More recently, several extensions of temporal logics and logics of coalitional ability that are

capable of expressing resource bounds have been proposed in the literature, e.g. [2, 7, 8, 10, 11].

All of these papers consider only the model-checking problem, and some of the logics allow both

consumption and production of resources by actions. There are many different proposals for the

syntax and semantics of resource logics. In [8] several versions are given, e.g. considering resource

bounds both on the coalition A and the rest of the agents in the system, considering a fixed resource

endowment of A in the initial state which affects their endowment after executing some actions, etc.

In [10, 11], a different syntax and semantics are considered, also involving resource endowment of

the whole system when evaluating a statement concerning a group of agents A. As observed in [8],

subtle differences in truth conditions for resource logics result in the difference between decidability
and undecidability of the model-checking problem. In [8], undecidability of the model-checking problem for several versions of the logics is proved. The only decidable cases considered in [8] are an extension of Computation Tree Logic (CTL) [9] with resources (essentially one-agent ATL) and the version where on every path only a fixed finite amount of resources can be produced. Similarly, [10] gives a logic PRB-ATL (Priced Resource-Bounded ATL) with a decidable model-checking problem where the total amount of resources in the system has a fixed bound. The model-checking algorithm for PRB-ATL runs in time polynomial in the size of the model and exponential in the number of resources and the resource bound on the system. In [11], an EXPTIME lower bound in the number of resources is shown. Recently, it has also been shown that if a zero-cost action is always available, the model-checking problem for RB-ATL with both production and consumption of resources is decidable; however, it is EXPSPACE-hard [2].

6 Conclusions
We have provided a complete and sound axiomatization of RB-ATL, a logic which extends ATL with resource bounds. The resulting logic can express interesting properties of coalitions of agents involving resource limitations. For example, it can express that a coalition can maintain the system in a $\varphi$-state indefinitely given a finite amount of resources (this essentially means that after a while $\varphi$ can be maintained for free). We have also presented a model-checking algorithm for RB-ATL, which runs in time polynomial in the size of the model and the formula (assuming that resource bounds are encoded in unary) and exponential in the number of resources.

The semantics for RB-ATL presented in this article, in particular the assumption that actions only consume but never produce resources, is motivated by verifying resource requirements for systems of agents where resources of interest are time, memory, bandwidth, etc., which cannot be generated by agents. In future work, we plan to study axiomatizations of variants of RB-ATL where actions can have a negative cost, such as in [2, 7].

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