Information Bottleneck on General Alphabets

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Abstract—We prove rigorously a source coding theorem that can probably be considered folklore, a generalization to arbitrary alphabets of a problem motivated by the Information Bottleneck method. For general random variables \((Y, X)\), we show that there exists a random variable \(U\) such that the Markov chain \(Y \rightarrow X \rightarrow U\) holds, \(I(U; X) \leq R\) and \(I(U; Y) \geq S\). The proof relies on the well-established discrete case and showcases a technique for lifting discrete coding theorems to arbitrary alphabets.

I. INTRODUCTION

Since its inception [1], the Information Bottleneck (IB) method became a widely applied tool, especially in the context of machine learning problems. It has been successfully applied to various problems in machine learning [2], computer vision [3], and communications [4], [5], [6]. Furthermore, it is a valuable tool for channel output compression in a communication system [7], [8].

In the underlying information-theoretic problem, we define a pair \((S, R) \in \mathbb{R}^2\) to be achievable for the two arbitrary random sources \((Y, X)\), if there exists a function \(f\) with rate limited range \(\frac{1}{n}\log|\mathcal{F}| \leq R\) and \(I(Y; f(X)) \geq nS\), where \((Y, X)\) are independent and identically distributed (i.i.d.) copies of \((Y, X)\).

While this Shannon-theoretic problem and variants thereof were also considered (e.g., [9], [10]), a large part of the literature is aimed at studying the IB function

\[ S_{IB}(R) = \sup_{U : I(U; X) \leq R} I(U; Y) = \sup_{U : I(U; X) \leq R} \frac{1}{n} I(Y; f(X)) \geq S. \]

in different contexts. In particular, several works (e.g., [1], [2], [11], [12], [13]) intend to compute a probability distribution that achieves the supremum in (1). The resulting distribution is then used as a building block in numerical algorithms, e.g., for document clustering [2] or dimensionality reduction [11].

In the discrete case, \(S_{IB}(R)\) is equal to the maximum of all \(S\) such that \((S, R)\) is in the achievable region (closure of the set of all achievable pairs). This statement has been re-proven many times in different contexts [14], [10], [15], [16]. In this note, we prove a theorem, which can probably be considered folklore, extending this result from discrete to arbitrary random variables. Formally speaking, using the definitions in [17], we prove that a pair \((S, R)\) is in the achievable region of an arbitrary source \((Y, X)\) if and only if, for every \(\epsilon > 0\), there exists a random variable \(U\) with \(Y \rightarrow X \rightarrow U\), \(I(X; U) \leq R + \epsilon\), and \(I(Y; U) \geq S - \epsilon\). This provides a single-letter solution to the information-theoretic problem behind the information bottleneck method for arbitrary random sources and in particular it shows, that the information bottleneck for Gaussian random variables [11] is indeed the solution to a Shannon-theoretic problem.

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The proof relies on the discrete case. Thus, the techniques employed could be useful for lifting other discrete coding theorems to the case of arbitrary alphabets.

II. MAIN RESULT

Let \(Y\) and \(X\) be random variables with arbitrary alphabets \(\mathcal{S}_Y\) and \(\mathcal{S}_X\), respectively. The bold-faced random vectors \(Y\) and \(X\) are \(n\) i.i.d. copies of \(Y\) and \(X\), respectively. We then have the following definitions.

Definition 1. A pair \((S, R) \in \mathbb{R}^2\) is achievable if for some \(n \in \mathbb{N}\) there exists a measurable function \(f : \mathcal{S}_X^n \rightarrow \mathcal{M}\) for some finite set \(\mathcal{M}\) with bounded cardinality \(\frac{1}{n}\log|\mathcal{M}| \leq R\) and

\[ \frac{1}{n} I(Y; f(X)) \geq S. \]

The set of all achievable pairs is denoted \(\mathcal{R}_{IB} \subseteq \mathbb{R}^2\).

Definition 2. A pair \((S, R) \in \mathbb{R}^2\) is IB-achievable if there exists an additional random variable \(U\) with arbitrary alphabet \(\mathcal{S}_U\), satisfying \(Y \rightarrow X \rightarrow U\) and

\[ R \geq I(X; U), \]

\[ S \leq I(Y; U). \]

The set of all IB-achievable pairs is denoted \(\mathcal{R}_{IB} \subseteq \mathbb{R}^2\).

In what follows, we will prove the following theorem.

Theorem 3. The equality \(\mathcal{R}_{IB} = \mathcal{R}\) holds.

III. PRELIMINARIES

When introducing a function, we implicitly assume it to be measurable w.r.t. the appropriate \(\sigma\)-algebras. The \(\sigma\)-algebra associated with a finite set is its power set and the \(\sigma\)-algebra associated with \(\mathbb{R}\) is the Borel \(\sigma\)-algebra. The symbol \(\mathcal{O}\) is used for the empty set and for a constant random variable. When there is no possibility for confusion, we will not distinguish between a single-element set and its element, e.g., we write \(x\) instead of \(\{x\}\) and \(1_x\) for the indicator function of \(\{x\}\). We use \(A \triangle B := (A \setminus B) \cup (B \setminus A)\) to denote the symmetric set difference.

Let \((\Omega, \Sigma, \mu)\) be a probability space. A random variable \(X : \Omega \rightarrow \mathcal{S}_X\) takes values in the measurable space \((\mathcal{S}_X, \mathcal{A}_X)\). The push-forward probability measure \(\mu_X : \mathcal{A}_X \rightarrow [0, 1]\) is defined by \(\mu_X(A) = \mu(X^{-1}(A))\) for all \(A \in \mathcal{A}_X\). We will state most results in terms of push-forward measures and usually ignore the background probability space. When multiple random variables are defined, we implicitly assume the push-forward measures to be consistent in the sense that, e.g., \(\mu_X(A) = \mu_X(X^{-1}(A))\) for all \(A \in \mathcal{A}_X\).

For \(n \in \mathbb{N}\) let \(\Omega^n\) denote the \(n\)-fold Cartesian product of \((\Omega, \Sigma, \mu)\). A bold-faced random vector, e.g., \(\mathbf{X}\), defined on \(\Omega^n\), is an \(n\)-fold copy of \(X\), i.e., \(\mathbf{X} = X^n\). Accordingly, the corresponding push-forward measure, e.g., \(\mu_X\) is the \(n\)-fold product measure.

For a random variable \(X\) let \(a_X, b_X\), and \(c_X\) denote arbitrary functions on \(\mathcal{S}_X\), each with finite range. We will use the symbol \(\mathcal{M}_X\) to denote the range of \(a_X\), i.e., \(a_X : \mathcal{S}_X \rightarrow \mathcal{M}_X\).
The conditional probability \( (iv) \) defined as \([17, \text{Lem. 5.5.7}]\) conditional entropy \( D(\cdot) \) is given by \([17, \text{Sec. 1.5}]\), there even exists a \([18, \text{Thm. 8.12}]\). Furthermore, if \( \mathbb{X} \rightarrow \mathbb{M}_X \) and \( \mathbb{Y} \rightarrow \mathbb{M}_Y \) be functions with finite range. Pick two arbitrary sets \( A \subseteq \mathbb{M}_X, B \subseteq \mathbb{M}_Y \) and we obtain \( \mu \)-a.e.

\[
\begin{align*}
\kappa_{\mathbb{X}Y}(\alpha_X | \alpha_Y) &= \kappa_{\mathbb{X}Y}(\cdot ; B) \quad (9) \\
&= \kappa_{\mathbb{X}Y}(\cdot ; B \mid \mathbb{Y}) \quad (10)
\end{align*}
\]

where \((9)\) and \((11)\) follow from part 3 of Def. 5. This proves that \( \mu \)-a.e. the equality of measures \( \kappa_{\mathbb{X}Y}(\alpha_X | \alpha_Y) = \kappa_{\mathbb{X}Y}(\alpha_X | \alpha_Y \mid \mathbb{Y}) \) holds. By the properties of Kullback-Leibler divergence \([17, \text{Thm. 2.3.1}]\) we have \( I(\mathbb{X}; \mathbb{Y}) = 0 \) due to \((6)\).

On the other hand, assume \( I(\mathbb{X}; \mathbb{Y}) = 0 \) and choose arbitrary sets \( A \subseteq \mathbb{A}_X \) and \( B \subseteq \mathbb{A}_Y \). We define \( a_X := \mathbb{1}_A, a_Y := \mathbb{1}_B, \hat{X} := a_X(\mathbb{X}), \) and \( \hat{Y} := a_Y(\mathbb{Y}) \). By \((6)\) we have \( D(\hat{X} | \hat{Y}) = D(\hat{X} | \hat{Y} \mid \mathbb{Y}) = 0 \) for \( \mu \)-a.e. \( \omega \in \mathbb{Y} \), which is equivalent to the equality \( \mu \)-a.e. of the measures \( \kappa_{\mathbb{X}Y} = \kappa_{\mathbb{X}Y} \mid \mathbb{Y} \). We obtain \( \mu \)-a.e.,

\[
\begin{align*}
P(\{X^{-1}(A) \cap Y^{-1}(B)\} | \mathbb{Y}) &= \kappa_{\mathbb{X}Y}(\cdot ; 1 \times 1) \\
&= \kappa_{\mathbb{X}Y}(\cdot ; 1) \kappa_{\mathbb{X}Y}(\cdot ; 1) \\
&= P(\{X^{-1}(A)\} | \mathbb{Z}) P(\{Y^{-1}(B)\} | \mathbb{Y}).
\end{align*}
\]

(ii): See \([17, \text{Lem. 5.5.6}]\).

(iii): See \([17, \text{Lem. 5.5.7}]\).

(iv): Using Prop. \((i)\) we have \( I(\mathbb{X}; \mathbb{Z}) = 0 \) and by Prop. \((iii)\) it follows that

\[
I(\mathbb{X}; \mathbb{Z}) \leq I(\mathbb{X}; \mathbb{Y}) + I(\mathbb{Z}; \mathbb{Y}) = I(\mathbb{X}; \mathbb{Y}).
\]

Occasionally we will interpret a probability measure on a finite space \( \mathcal{M} \) as a vector in \([0,1]^{|\mathcal{M}|}\), equipped with the Borel \( \sigma \)-algebra. We will use the \( L_{\infty} \)-distance on this space.

**Definition 9.** For two probability measures \( \mu \) and \( \nu \) on a finite space \( \mathcal{M} \), their distance is defined as the \( L_{\infty} \)-distance \( d(\mu, \nu) := \max_{m \in \mathcal{M}} |\mu(m) - \nu(m)| \). The diameter of \( A \subseteq \{0,1\}^{|\mathcal{M}|} \) is defined as \( \text{diam}(A) = \sup_{\mu, \nu \in A} d(\mu, \nu) \).

**Lemma 10.** \([19, \text{Lem. 2.7.}]\). For two probability measures \( \mu \) and \( \nu \) on a finite space \( \mathcal{M} \) with \( d(\mu, \nu) \leq \varepsilon < \frac{1}{2} \), the inequality \( |H(\mu) - H(\nu)| \leq -\varepsilon |\mathcal{M}| \log \varepsilon \) holds.

**IV. PROOF OF \( R_{IB} \subseteq \mathcal{F} \)**

For finite spaces \( \mathcal{X}, \mathcal{X}, \text{and } \mathcal{B} \), the statement \( \mathcal{R}_{IB} \subseteq \mathcal{F} \) is well known, cf., \([9, \text{Sec. IV}]\), \([10, \text{Sec. III.F}]\). We restate it in the form of the following lemma.

**Lemma 11.** For random variables \( \mathbb{Y}, \mathbb{X}, \text{and } \mathcal{U} \) with finite \( \mathcal{X}, \mathcal{X}, \text{and } \mathcal{U} \), assume that \( \mathbb{Y} \rightarrow \mathcal{X} \rightarrow \mathcal{U} \) holds. Then, for any \( \varepsilon > 0 \), there exists \( n \in \mathbb{N} \) and a function \( f: \mathcal{X}_n \rightarrow \mathcal{M} \) with \( \frac{1}{n} \log |\mathcal{M}| \leq I(\mathbb{X}; \mathcal{U}) + \varepsilon \) such that \( \frac{1}{n} I(\mathbb{Y}; f(\mathbb{X})) \geq I(\mathbb{Y}; \mathcal{U}) - \varepsilon \).

In a first step, we will utilize Lem. 11 to show \( \mathcal{R}_{IB} \subseteq \mathcal{F} \) for an arbitrary alphabet \( \mathcal{X}_n \), i.e., we wish to prove the following Proposition 12, lifting the restriction \( |\mathcal{X}| < \infty \).

**Proposition 12.** For random variables \( \mathbb{Y}, \mathbb{X}, \text{and } \mathcal{U} \) with finite \( \mathcal{X}, \mathcal{X}, \text{and } \mathcal{U} \), assume that \( \mathbb{Y} \rightarrow \mathcal{X} \rightarrow \mathcal{U} \) holds. Then, for any
as the according projections. We also use \( \mu \) probability space random variable probability measures on and find a finite, measurable partition Proof of Proposition 12.

\[ I(\hat{Y}; U) = \int \kappa_{YXU} (\cdot; y \times u) \, d\mu_{XU} \]

\[ \geq \mu_{YXU}(y \times u) \mu_{XU}(u) \]

Lemma 13.

\[ \mathcal{R}_{IB} \subseteq \overline{\mathcal{R}}. \]

**Remark 2.** Considering that both definitions of achievability (Defs. 1 and 2) only rely on the notion of mutual information, one may assume that Def. 6 can be used to directly infer Proposition 12 from Lem. 11. However, this is not the case. For an arbitrary discretization \( \alpha_X(Y) \) of \( X \), we do have \( I(\alpha_X(Y); U) \leq I(Y; U) \). However, the Markov chain \( Y \rightarrow X \rightarrow U \) does not hold in general. To circumvent this problem, we will use a discrete random variable \( \tilde{X} = g(X) \) with an appropriate quantizer \( g \) and construct a new random variable \( U \), satisfying the Markov chain \( Y \rightarrow \tilde{X} \rightarrow U \) such that \( I(Y; U) \) is close to \( I(Y; U) \). Fig. 1a illustrates this strategy. We choose the quantizer \( g \) based on the conditional probability distribution of \( U \) given \( X \), i.e., quantization based on \( \kappa_{YU|X} \) using \( L_\infty \)-distance (c.f. Def. 9). Subsequently, we will use that, by Lem. 10, a small \( L_\infty \)-distance guarantees a small gap in terms of information measures.

**Proof of Proposition 12.** Let \( \mu_{XU} \) be a probability measure on \( \Omega := S_Y \times S_X \times S_U \), such that \( Y \rightarrow X \rightarrow U \) holds. Fix \( 0 < \delta \leq \frac{1}{2} \) and find a finite, measurable partition \((P_i)_{i \in I}\) of the space of probability measures on \( S_U \) such that for every \( i \in I \) we have \( \text{diam}(P_i) \leq \delta \) and fix some \( \nu_i \in P_i \) for every \( i \in I \). Define the random variable \( \tilde{X} : \Omega \rightarrow I \) as \( \tilde{X} = i \) if \( \kappa_{YU|X} \in P_i \). The random variable \( \tilde{X} \) is \( \sigma(X) \)-measurable (see Appendix A). We can therefore find a measurable function \( g \) such that \( \tilde{X} = g(X) \) by the factorization lemma [18, Corollary 1.97]. Define the new probability space \( \Omega \times \times \in \mathcal{E} \), equipped with the probability measure \( \mu_{YXU} := \mu_{XU} \times \times \mu_{U} \). Slightly abusing notation, we define the random variables \( Y, X, U, \) and \( \tilde{U} \) for every \( i \in I \) as the according projections. We also use \( \tilde{X} = g(X) \) and define the random variable \( \tilde{U} = \tilde{U}_X \). From this construction we have \( \kappa_{YXU} \)-a.e. the equality of measures \( \kappa_{\tilde{X}U|X} = \kappa_{\tilde{U}|X} = \kappa_{X} \), as well as \( Y \rightarrow \tilde{X} \rightarrow \tilde{U} \) and \( Y \rightarrow X \rightarrow U \). This is proven in the extended version [20]. Therefore, we have \( \kappa_{YXU} \)-a.e.

\[ d(\kappa_{\tilde{X}U|X}, \kappa_{U|X}) \leq \delta. \]

Thus, we obtain

\[ I(\hat{Y}; U) \geq I(Y; U) - \delta \]

\[ \geq I(Y; U) - \frac{1}{n} \log |\mathcal{R}| \]
functions $a_n$, $b_n$ such that $I(a_n(X); a_n(U)) \geq I(X; U) - \epsilon$, and select functions $b_n$, $b_0$ such that $I(b_n(Y); b_0(U)) \geq I(Y; U) - \epsilon$ (cf. (7)). Using $\hat{U} := (a_n(U), b_n(U))$ and $\hat{Y} := b_0(Y)$, we have

$$0 = I(Y; U | X) = \sup_{\phi_X, \psi_U} I(\phi_{cY}(Y); \psi_{cU}(U) | X) \geq I(\hat{Y}; \hat{U} | X) \geq 0 \quad (38)$$

as well as

$$I(X; U) = \sup_{\phi_X, \psi_U} I(\phi_{cX}(X); \psi_{cU}(U)) \quad (39)$$

$$\geq \sup_{\phi_X} I(\phi_{cX}(X); \hat{U}) = I(\hat{X}; \hat{U}), \text{ and} \quad (40)$$

$$I(Y; U) - \epsilon \leq I(\hat{Y}(Y); b_0(U)) \leq I(\hat{Y}; \hat{U}). \quad (41)$$

We apply Proposition 12, substituting $\hat{U} \to U$ and $\hat{Y} \to Y$. Proposition 12 guarantees the existence of a function $f : S'_X \to \mathcal{M}$ with $\frac{1}{n} \log |\mathcal{M}| \leq I(X; \hat{U}) + \epsilon \leq I(X; U) + \epsilon \leq R + \epsilon$ and

$$\frac{1}{n} I(Y; f(X)) = \frac{1}{n} \sup_{\phi_Y} I(\phi_{cY} \circ Y; f(X)) \leq \frac{1}{n} I(\hat{Y}; f(X)) \leq I(\hat{Y}; \hat{U}) - \epsilon \leq I(Y; U) - 2\epsilon \geq S - 2\epsilon. \quad (43)$$

Thus, $(S - 2\epsilon, R - \epsilon) \in \mathcal{R}$ and therefore $(S, R) \in \overline{\mathcal{R}}$.

V. PROOF OF $\mathcal{R} \subseteq \overline{\mathcal{R}}$

We start with the well-known result $\mathcal{R}_{1B} \subseteq \overline{\mathcal{R}}$ for finite spaces $S_Y$, $S_X$, and $S_U$, cf., [9, Sec. IV], [10, Sec. III.F]. The statement is rephrased in the following lemma.

**Lemma 14.** Assume that the spaces $S_Y$ and $S_X$ are both finite and $\mu_{XY}$ is fixed. For some $n \in \mathbb{N}$, let $f : S_n^X \to \mathcal{M}$ be a function with $|\mathcal{M}| < \infty$. Then there exists a probability measure $\mu_{XY}$, extending $\mu_{XY}$, such that $\hat{S}_U$ is finite, $Y \to X \to U$, and

$$I(X; U) \leq \frac{1}{n} \log |\mathcal{M}|, \quad (44)$$

$$I(Y; U) \geq \frac{1}{n} I(Y; f(X)). \quad (45)$$

We can slightly strengthen Lemma 14.

**Corollary 15.** Assume that, in the setting of Lem. 14, we are given $\mu_{XY}$ on $S_Y \times S_X \times S_U$, extending $\mu_{XY}$, where $S_U$ is arbitrary, not necessarily finite. Then there exists a probability measure $\mu_{XY}$, extending $\mu_{XY}$, such that $\hat{S}_U$ is finite and $Y \to X \to U$, (45), and (46) hold.

**Proof.** Apply Lem. 14 to obtain $\mu_{XY}$ on $S_Y \times S_X \times S_U$ satisfying (45), (46), and $Y \to X \to U$. We define $\mu_{XY}$ by

$$\mu_{XY}(A \times y \times x \times u) = \frac{\mu_{XY}(A \times y \times x)}{\mu_{XY}(y \times x)} \mu_{XY}(y \times x \times u) \quad (47)$$

for any $(y, x, u) \in S_Y \times S_X \times S_U$ and $A \in \mathcal{A}_2$. Pick arbitrary $A \in \mathcal{A}_2$, $y \in S_Y$, and $u \in S_U$. The Markov chain $Z \to Y \to X \to U$ now follows as the events $Z^{-1}(A) \cap Y^{-1}(y)$ and $U^{-1}(u)$ are independent given $X^{-1}(x)$ for any $x \in S_X$ (cf. Rumk. 1).}

Again, we proceed by extending Cor. 15, lifting the restriction that $S_X$ is finite and obtain the following proposition.

**Proposition 16.** Given a probability measure $\mu_{XY}$ as in Cor. 15, assume that $|S_Y| < \infty$. For some $n \in \mathbb{N}$, let $f : S_n^Y \to \mathcal{M}$ be a function with $|\mathcal{M}| < \infty$. Then, for any $\epsilon > 0$, there exists a probability measure $\mu_{XY}$, extending $\mu_{XY}$ with $Z \to X \to U$ and

$$I(X; U) \leq \frac{1}{n} \log |\mathcal{M}| \quad (48)$$

$$I(Y; U) \geq \frac{1}{n} I(Y; f(X)) - \epsilon. \quad (49)$$

**Remark 3.** In contrast to Proposition 12, Proposition 16 could be proved by the usual single-letterization + time-sharing strategy, by showing that the necessary Markov chains hold. However, we will rely on the discrete case (Lem. 14) and showcase a technique to lift it to general alphabets.

**Remark 4.** In the proof of Proposition 16, we face a similar problem as outlined in Rumk. 2. We need to construct a function $g(X)$ of a “per-letter” quantization $X := a_2^X$, that is close to $f(X)$ in distribution. Fig. 1b provides a sketch.

**Proof of Proposition 16.** We can partition $S_n^Y = \bigcup_{m=1}^{\infty} S_n^m$ into finitely many measurable, mutually disjoint sets $Q_m := f^{-1}(m), m \in \mathcal{M}$. We want to approximate the sets $Q_m$ by a finite union of rectangles in the semiring $[18, \text{Def. 1.9}] \Xi := \{B : B = \bigcup_{i=1}^{n} B_i, B_i \in \mathcal{A}_X\}$. We choose $\delta > 0$, which will be specified later. According to [18, Thm. 1.65(ii)], we obtain $B(m) := \bigcup_{n=1}^{\infty} B_n(m)$ for each $m \in \mathcal{M}$, where $B_n(m) \in \Xi$ are mutually disjoint sets, satisfying $\mu_X(B(m) \cap Q_m) \leq \delta$. Indeed, we obtain $\mu_X$ by finding a measurable partition of $S_X$ that is finer than $(B_n(m), (B_n(m))')$ for all $i, k, m$. For fixed $m \in \mathcal{M}$,

$$Q_m \subseteq Q_m \cup (B(m) \setminus B_n(m)) \quad (50)$$

$$\subseteq (B(m) \setminus B_n(m)) \cup \bigcup_{m' \neq m} Q_m \cap B(m') \quad (51)$$

$$\subseteq (B(m) \setminus B_n(m)) \cup (Q_m \setminus B(m)) \cup \bigcup_{m' \neq m} B_n(m') \cap Q_m \quad (52)$$

$$\subseteq (B(m) \setminus B_n(m)) \cup \bigcup_{m' \neq m} B_n(m') \cap Q_m' \quad (53)$$

where we used the fact that $Q_m \cap Q_m' = \emptyset$ for $m \neq m'$ in (52).

Using $X := a_2^X$, we obtain for any $y \in S_n^Y$

$$\mu_{XY}(y \times m) = \mu_{XY}(y \times Q_m) \quad (54)$$

$$\leq \mu_{XY}(y \times (B(m) \setminus B_n(m)) + \bigcup_{m' \neq m} \mu_{XY}(B_n(m) \cap Q_m') \quad (55)$$

$$\mu_{XY}(y \times m) + |\mathcal{M}| \delta. \quad (56)$$

On the other hand, we have

$$\mu_{XY}(y \times m) = \mu_{XY}(y) - \sum_{m' \neq m} \mu_{XY}(y \times m') \quad (57)$$

$$\leq \mu_{XY}(y) - \sum_{m' \neq m} (\mu_{XY}(y \times m') + |\mathcal{M}| \delta) \quad (58)$$

$$\mu_{XY}(y \times m) - |\mathcal{M}| \delta. \quad (59)$$

We thus obtain $\mu_{XY}(y \times m') \leq |\mathcal{M}| \delta$. This also implies $|\mathcal{S}_Y| |\mathcal{M}| \delta \leq \frac{1}{2}$ and apply Cor. 15 substituting $\hat{X} \to X$, $XZ \to Z$, and the function $g \to f$. This yields a random variable $U$ with $XZ \to U$, $I(X; U) \leq \frac{1}{n} \log |\mathcal{M}|$, and $I(Y; U) \geq \frac{1}{n} I(Y; g(\hat{X})).$ (60)
We also obtain $ZY \rightarrow X \rightarrow U$ due to

$$\begin{align*}
0 &= I(X; Y; U | X) \\
&= I(X; Z; U | X) - I(X; U | X) \\
&\geq I(X; Z; U | X) - I(U; X) \\
&= I(Z; U | X) \\
&\geq 0,
\end{align*}$$

(61)

where (61) follows from $XZ \rightarrow Y \rightarrow U$ using Prop. (i) of Lem. 8, (62) and (64) follow from Prop. (iii) of Lem. 8, (63) is a consequence of Def. 6, and we used Prop. (i) of Lem. 8 in (65). This also immediately implies $0 = I(X; U | X)$ and hence

$$\begin{align*}
\frac{1}{n} \log |M| &\geq I(X; U | X) = I(X; U | X) + I(X; U | X) \\
&\geq I(X; U | X),
\end{align*}$$

(66)

where we used Prop. (iii) of Lem. 8 in (67). We also have

$$\begin{align*}
I(Y; U) &\geq \frac{1}{n} I(Y; g(X)) \\
&\geq \frac{1}{n} \left( H(Y) + H(g(X)) - H(Y | g(X)) \right) \\
&\geq \frac{1}{n} \left[ I(Y; f(X)) + \frac{1}{n} |S| n^{-3} \delta \log (|M| \delta) \right] \\
&\quad + \frac{1}{n} |S| n^{-3} \delta \log (|S| n^{-3} |M| \delta^2) \\
&\geq \frac{1}{n} \left[ I(Y; f(X)) + \frac{2}{n} |S| n^{-3} \delta \log (|M| \delta^2) \right]
\end{align*}$$

(68)

where we used Lem. 10 in (70). Select $\delta$ such that $\varepsilon \geq -\frac{2}{n} |S| n^{-3} \delta \log (|M| \delta^2)$. ■

We can now finish the proof by showing the following lemma.

**Lemma 17.** $R \subseteq \overline{R}_{IB}$.

**Proof.** Assume $(S, R) \in R$ and choose $n \in N$ and $f$, satisfying $\frac{1}{n} \log |M| \leq R$ and (2). Choose any $\varepsilon > 0$ and find $\sigma$ such that

$$I(\sigma(Y); f(X)) \geq I(Y; f(X)) - \varepsilon \geq nS - \varepsilon.$$

(72)

This is possible by applying [17, Lem. 5.2.2] with the algebra that is generated by the rectangles (cf. the paragraph above [17, Lem. 5.5.1]). We apply Proposition 16, substituting $\sigma(Y) \rightarrow Y$ and $Y \rightarrow Z$. For arbitrary $\varepsilon > 0$, Proposition 16 provides $U$ with $Y_{AV}(Y) \rightarrow X \rightarrow U$ (i.e., $Y \rightarrow X \rightarrow U$) and

$$\begin{align*}
I(X; U) &\leq \frac{1}{n} \log |M| \leq R \\
I(Y; U) &\geq I(\sigma(Y); U) \\
&\geq \frac{1}{n} I(\sigma(Y); f(X)) - \varepsilon \geq S - 2\varepsilon
\end{align*}$$

(73)

Hence, $(S - 2\varepsilon, R) \in \overline{R}_{IB}$ and consequently $(S, R) \in \overline{R}_{IB}$. ■

**Appendix**

A. $\tilde{X}$ is $\sigma(X)$-measurable

For $u \in S_\theta$ consider the $\sigma(X)$-measurable function $h_u := \kappa_{\Theta} (\cdot; u)$ on $[0, 1]$. We obtain the vector valued function $h := (h_u)_{u \in S_\theta}$ on $[0, 1]^{|S_\theta|}$. This function $h$ is $\sigma(X)$-measurable as every component is $\sigma(X)$-measurable. Thus, we have $\tilde{X}^{-1}(i) = h^{-1}(P_i) \in \sigma(X)$.

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