Box-total dual integrality and edge-connectivity

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Abstract
Given a graph $G = (V, E)$ and an integer $k \geq 1$, the graph $H = (V, F)$, where $F$ is a family of elements (with repetitions allowed) of $E$, is a $k$-edge-connected spanning subgraph of $G$ if $H$ cannot be disconnected by deleting any $k - 1$ elements of $F$. The convex hull of incidence vectors of the $k$-edge-connected subgraphs of a graph $G$ forms the $k$-edge-connected subgraph polyhedron of $G$. We prove that this polyhedron is box-totally dual integral if and only if $G$ is series-parallel. In this case, we also provide an integer box-totally dual integral system describing this polyhedron.

Keywords Box-total dual integrality · $k$-edge connected subgraph · Polyhedron · Series-parallel graph

Mathematics Subject Classification 90C27 · 90C05 · 90C57

1 Introduction

Totally dual integral systems, introduced in the late 70’s, are strongly connected to min–max relations in combinatorial optimization [34]. A rational system of linear
inequalities $Ax \geq b$ is totally dual integral (TDI) if the maximization problem in the linear programming duality

$$
\min\{c^T x : Ax \geq b\} = \max\{b^T y : A^T y = c, y \geq 0\}
$$

admits an integer optimal solution for each integer vector $c$ such that the optimum is finite. Every rational polyhedron can be described by a TDI system [28]. For instance, the polyhedron $\{x : Ax \geq b\}$ can be described by TDI systems of the form $\frac{1}{q} Ax \geq \frac{1}{q} b$ for certain positive $q$. However, a polyhedron is integer if and only if it can be described by a TDI system with only integer coefficients [23,28]. Integer TDI systems yield min–max results that may have combinatorial interpretation.

A stronger property is box-total dual integrality: a system $Ax \geq b$ is box-totally dual integral (box-TDI) if $Ax \geq b$, $\ell \leq x \leq u$ is TDI for all rational vectors $\ell$ and $u$ (possibly with infinite components). General properties of such systems can be found in Cook [12] and Chapter 22.4 of Schrijver [34]. Note that, although every rational polyhedron can be described by a TDI system, not every polyhedron can be described by a box-TDI system. A polyhedron which can be described by a box-TDI system is called a box-TDI polyhedron. As proved by Cook [12], every TDI system describing such a polyhedron is actually box-TDI.

Recently, several new box-TDI systems have been exhibited. Chen et al. [6] characterized box-Mengerian matroid ports. Ding et al. [18] characterized the graphs for which the Edmonds’ system defining the matching polytope [21] is box-TDI. Ding et al. [19] exhibited new subclasses of box-perfect graphs. Cornaz et al. [14] provided several box-TDI systems in series–parallel graphs. Barbato et al. [3] gave the minimal box-TDI system with integer coefficients for the flow cone for series–parallel graphs. For these graphs, Chen et al. [7] provided a box-TDI system describing the 2-edge-connected spanning subgraph polyhedron.

In this paper, we are interested in integrality properties of systems related to $k$-edge-connected spanning subgraphs. A $k$-edge-connected spanning subgraph of a graph $G = (V, E)$ is a graph $H = (V, F)$, with $F$ being a collection of elements of $E$ where each element can appear several times, that remains connected after the removal of any $k - 1$ edges.

These objects model a kind of failure resistance of telecommunication networks. More precisely, they represent networks which remain connected when $k - 1$ links fail. The underlying network design problem is the $k$-edge-connected spanning subgraph problem ($k$-ECSPP): given a graph $G$ and positive edge costs, find a $k$-edge-connected spanning subgraph of $G$ of minimum cost. Special cases of this problem are related to classical combinatorial optimization problems. The 2-ECSPP is a well-studied relaxation of the traveling salesman problem [24] and the 1-ECSPP is nothing but the well-known minimum spanning tree problem. While this latter is polynomial-time solvable, the $k$-ECSPP is NP-hard for every fixed $k \geq 2$ [27].

Different algorithms have been devised in order to deal with the $k$-ECSPP, such as branch-and-cut procedures [4,15], approximation algorithms [8,26], cutting plane algorithms [30], and heuristics [11]. In [36], Winter introduced a linear-time algorithm
solving the 2-ECSSP on series–parallel graphs. Most of these algorithms rely on polyhedral considerations.

Given a graph $G = (V, E)$, the convex hull of incidence vectors of all the families of $E$ inducing a $k$-edge-connected spanning subgraph of $G$ forms a polyhedron, hereafter called the $k$-edge-connected spanning subgraph polyhedron of $G$ and denoted by $P_k(G)$. Cornuéjols et al. [16] gave a system describing $P_2(G)$ when $G$ is series–parallel. Vandenbussche and Nemhauser [35] characterized in terms of forbidden minors the graphs for which this system describes $P_2(G)$. Chopra [10] described $P_k(G)$ for outerplanar graphs when $k$ is odd. Didi Biha and Mahjoub [17] extended these results to series–parallel graphs for all $k \geq 2$. By a result of Baïou et al. [1], the inequalities in these descriptions can be separated in polynomial time, which implies that the $k$-ECSSP is solvable in polynomial time for series–parallel graphs.

When studying $k$-edge-connected spanning subgraphs of a graph $G$, we can add the constraint that each edge of $G$ can be taken at most once. We denote the corresponding polyhedron by $Q_k(G)$. Barahona and Mahjoub [2] described $Q_2(G)$ for Halin graphs. Further polyhedral results for the case $k = 2$ have been obtained by Boyd and Hao [5] and Mahjoub [32,33]. Grötschel and Monma [29] described several classes of facets of $Q_k(G)$. Moreover, Fonlupt and Mahjoub [25] extensively studied the extremal points of $Q_k(G)$ and characterized the class of graphs for which this polytope is described by cut inequalities and $0 \leq x \leq 1$.

The polyhedron $P_1(G)$ is known to be box-TDI for all graphs [31]. For series–parallel graphs, the system given in [16] describing $P_2(G)$ is not TDI. Chen et al. [7] showed that dividing it by 2 yields a TDI system for such graphs. Actually, they proved that this system is box-TDI if and only if the graph is series–parallel.

Contributions. Our starting point is the result of Chen et al. [7]. First, their result implies that $P_2(G)$ is a box-TDI polyhedron for series–parallel graphs. However, this leaves open the question of the box-TDIness of $P_2(G)$ for non series–parallel graphs. More generally, for which integers $k$ and graphs $G$ is $P_k(G)$ a box-TDI polyhedron?

We answer this question by proving that, for $k \geq 2$, $P_k(G)$ is a box-TDI polyhedron if and only if $G$ is series–parallel. Note that this work is one of the first ones that proves the box-TDIness of a polyhedron without giving a box-TDI system describing it. Instead, our proof is based on the recent matricial characterization of box-TDI polyhedra given by Chervet et al. [9].

By [34, Theorem 22.6], there exists a TDI system with integer coefficients describing $P_k(G)$. For series–parallel graphs, the system provided by Chen et al. [7] has noninteger coefficients. Moreover, the system given by Didi Biha and Mahjoub [17] describing $P_k(G)$ when $k$ is even is not TDI. When $k \geq 2$ and $G$ is series–parallel, which combinatorial objects yield an integer TDI system describing $P_k(G)$?

We answer this question by exhibiting integer TDI systems based on multicuts. When $k$ is even, we use multicuts to provide an integer TDI system for $P_k(G)$ when $G$ is series–parallel. Our proof relies on the standard constructive characterization of series–parallel graphs. When $k$ is odd, we prove that the description of $P_k(G)$ given by Didi Biha and Mahjoub [17] based on multicuts is TDI if and only if the graph is series–parallel. For this case, our proof relies on new properties of the set of degree 2 vertices in simple series–parallel graphs stated in Lemma 2.3.
The box-totally dual integral characterization of $P_k(G)$ implies that these systems are actually box-TDI if and only if $G$ is series–parallel. By definition of box-TDIness, adding $x \leq 1$ to these systems yields box-TDI systems for $Q_k(G)$ for series–parallel graphs.

Outline. In Sect. 2, we give the definitions and preliminary results used throughout the paper. In Sect. 3, we prove that, for $k \geq 2$, $P_k(G)$ is a box-TDI polyhedron if and only if $G$ is series–parallel. By definition of box-TDIness, adding $x \leq 1$ to these systems yields box-TDI systems for $Q_k(G)$ for series–parallel graphs.

2 Definitions and preliminary results

This section is devoted to the definitions, notation, and preliminary results used throughout the paper.

2.1 Graphs and combinatorial objects

Given a set $E$, a family of $E$ is a collection of elements of $E$ where each element can appear multiple times. The incidence vector of a family $F$ of $E$ is the vector $\chi_F$ of $\mathbb{Z}_E^+$ such that $e$’s coordinate is the multiplicity of $e$ in $F$ for all $e$ in $E$. Since there is a bijection between families and their incidence vectors, we will often use the same terminology for both.

Given a graph $G = (V, E)$ and the incidence vector $z \in \mathbb{Z}_E^+$ of a family $F$ of $E$, $G(z)$ denotes the graph $(V, F)$.

Let $G = (V, E)$ be a loopless undirected graph. Two edges of $G$ are parallel if they share the same endpoints, and $G$ is simple if it does not have parallel edges. A graph is 2-connected if it cannot be disconnected by removing a vertex. The graph obtained from two disjoint graphs by identifying two vertices, one of each graph, is called a 1-sum. A 2-connected graph is trivial if it is composed of a single edge. We denote by $K_n$ the complete graph on $n$ vertices, that is the simple graph with $n$ vertices and one edge between each pair of vertices. Given an edge $e$ of $G$, we denote by $G \setminus e$ (respectively $G/e$) the graph obtained by removing (respectively contracting) the edge $e$, where contracting an edge $uv$ consists in removing it and identifying $u$ and $v$. Similarly, we denote by $G \setminus v$ the graph obtained from $G$ by removing the vertex $v$, and by $G[W]$ the graph induced by $W$, that is, the graph obtained by removing all vertices not in the vertex subset $W$. Given a vector $x \in \mathbb{R}^E$ and a subgraph $H$ of $G$, we denote by $x|_H$ the vector obtained by restricting $x$ to the components associated with the edges of $H$.

A subset of edges of $G$ is called a circuit if it induces a connected graph in which every vertex has degree 2. Given a subset $U$ of $V$, the cut $\delta(U)$ is the set of edges having exactly one endpoint in $U$. A bond is a minimal nonempty cut. Given a partition $\{V_1, \ldots, V_n\}$ of $V$, the set of edges having endpoints in two distinct $V_i$’s is called a multicut and is denoted by $\delta(V_1, \ldots, V_n)$. We denote respectively by $\mathcal{M}_G$ and $\mathcal{B}_G$ the
set of multicuts and the set of bonds of $G$. For every multicut $M$, there exists a unique partition $\{V_1, \ldots, V_{d_M}\}$ of vertices of $V$ such that $M = \delta(V_1, \ldots, V_{d_M})$, and $G[V_i]$ is connected for all $i = 1, \ldots, d_M$. We say that $d_M$ is the order of $M$ and $V_1, \ldots, V_{d_M}$ are the classes of $M$. Multicuts are characterized in terms of circuits, as stated in the following.

**Lemma 2.1** [13] A set of edges $M$ is a multicut if and only if $|M \cap C| \neq 1$ for all circuits $C$ of $G$.

We denote the symmetric difference of two sets $S$ and $T$ by $S \triangle T$. It is well-known that the symmetric difference of two cuts is a cut. Moreover, the following result holds.

**Observation 2.2** Let $G$ be a graph, $v$ be a degree 2 vertex of $G$, and $M$ be a multicut such that $|M \cap \delta(v)| = 1$. Then, $M \cup \delta(v)$ and $M \triangle \delta(v)$ are multicuts. Moreover, $d_{M \cup \delta(v)} = d_M + 1$, and $d_{M \triangle \delta(v)} = d_M$.

A graph is series–parallel if its nontrivial 2-connected components can be constructed from a circuit of length 2 by repeatedly adding edges parallel to an existing one, and subdividing edges, that is, replacing an edge by a path of length two. Equivalently, series–parallel graphs are those having no $K_4$-minor [20].

By construction, simple nontrivial 2-connected series–parallel graphs have at least one degree 2 vertex. Moreover, these vertices satisfy the following.

**Lemma 2.3** For a simple nontrivial 2-connected series–parallel graph, at least one of the following holds:

(i) Two degree 2 vertices are adjacent,
(ii) A degree 2 vertex belongs to a circuit of length 3,
(iii) Two degree 2 vertices belong to the same circuit of length 4.

**Proof** We proceed by induction, the base case is $K_3$ for which (i) holds.

Let $G$ be a simple 2-connected series–parallel graph. Since $G$ is simple, it can be built from a series–parallel graph $H$ by subdividing an edge $e$ into a path $f, g$. Let $v$ be the degree 2 vertex added with this operation. By the induction hypothesis, either $H$ is not simple, or one among (i), (ii), and (iii) holds for $H$. Hence, there are four cases.

**Case 1** $H$ is not simple. Since $G$ is simple, $e$ is parallel to exactly one edge $h$. Hence, $f, g, h$ is a circuit of $G$ length 3 containing $v$, thus (ii) holds for $G$.

**Case 2** (i) Holds for $H$. Then, it holds for $G$.

**Case 3** (ii) Holds for $H$. Let $C$ be a circuit of $H$ of length 3 containing a degree 2 vertex, say $w$. If $e \not\in C$, then (ii) holds for $G$. Otherwise, by subdividing $e$, we obtain a circuit of length 4 containing $v$ and $w$, and hence (iii) holds for $G$.

**Case 4** (iii) Holds for $H$. Let $C$ be a circuit of $H$ of length 4 containing two degree 2 vertices. If $e \not\in C$, then (iii) holds for $G$. Otherwise, by subdividing $e$, we obtain a circuit of length 5 containing three degree 2 vertices. Then, at least two of them are adjacent, and so (i) holds for $G$. \hfill \square
2.2 Box-total dual integrality

Let \( A \in \mathbb{R}^{m \times n} \) be a full-row rank matrix. This matrix is equimodular if all its \( m \times m \) non-zero determinants have the same absolute value. The matrix \( A \) is face-defining for a face \( F \) of a polyhedron \( P \subseteq \mathbb{R}^n \) if \( \text{aff}(F) = \{x \in \mathbb{R}^n : Ax = b\} \) for some \( b \in \mathbb{R}^m \), where \( \text{aff}(F) \) denotes the affine hull of \( F \). Such matrices are the face-defining matrices of \( P \).

**Theorem 2.4** [9, Theorem 1.4] Let \( P \) be a polyhedron. Then, the following statements are equivalent:

(i) \( P \) is box-TDI.
(ii) Every face-defining matrix of \( P \) is equimodular.
(iii) Each face of \( P \) has an equimodular face-defining matrix.

In Theorem 2.4, the equivalence of conditions (ii) and (iii) follows from the following observation.

**Observation 2.5** ([9, Observation 4.10]) Let \( F \) be a face of a polyhedron. If a face-defining matrix for \( F \) is equimodular, then so are all the face-defining matrices for \( F \).

We will also use the following.

**Observation 2.6** Let \( A \in \mathbb{R}^{I \times J} \) be a full row rank matrix and \( j \in J \). If \( A \) is equimodular, then so are following two matrices:

(i) \[
\begin{bmatrix}
A \\ \pm \chi j
\end{bmatrix}
\]
if it is full row-rank,

(ii) \[
\begin{bmatrix}
A \\ 0 \\ \pm \chi j \\ \pm 1
\end{bmatrix}
\]

**Observation 2.7** [9, Observation 4.11] Let \( P \subseteq \mathbb{R}^n \) be a polyhedron and let \( F = \{x \in P : Bx = b\} \) be a face of \( P \). If \( B \) has full-row rank and \( n - \dim(F) \) rows, then \( B \) is face-defining for \( F \).

2.3 The \( k \)-edge-connected spanning subgraph polyhedron

Note that \( P_k(G) \) is the dominant of the convex hull of incidence vectors of all the families of \( E \) containing at most \( k \) copies of each edge and inducing a \( k \)-edge-connected spanning subgraph of \( G \). Since the dominant of a polyhedron is a polyhedron, \( P_k(G) \) is a full-dimensional polyhedron even though it is the convex hull of an infinite number of points.

From now on, we assume that \( k \geq 2 \). Didi Biha and Mahjoub [17] gave a complete description of \( P_k(G) \) for all \( k \), when \( G \) is series–parallel.

**Theorem 2.8** ([17]) Let \( G \) be a series–parallel graph and \( h \) be a positive integer. Then, \( P_{2h}(G) \) is described by:

\[
\begin{align}
(1a) & \quad x(D) \geq 2h \quad \text{for all cuts } D \text{ of } G, \\
(1b) & \quad x \geq 0,
\end{align}
\]
and $P_{2h+1}(G)$ is described by:

$$
\begin{align*}
(2) \begin{cases} 
    x(M) &\geq (h + 1)d_M - 1 \quad \text{for all multicuts } M \text{ of } G, \\
    x &\geq 0. 
\end{cases}
\end{align*}
$$

Since the incidence vector of a multicut $\delta(V_1, \ldots, V_\ell)$ of order $\ell$ is the half-sum of the incidence vectors of the bonds $\delta(V_1), \ldots, \delta(V_\ell)$, we can deduce another description of $P_{2h}(G)$.

Corollary 2.9 Let $G$ be a series–parallel graph and $h$ be a positive integer. Then, $P_{2h}(G)$ is described by:

$$
\begin{align*}
(3) \begin{cases} 
    x(M) &\geq hd_M \quad \text{for all multicuts } M \text{ of } G, \\
    x &\geq 0. 
\end{cases}
\end{align*}
$$

We call constraints (2a) and (3a) partition constraints. A multicut $M$ is tight for a point of $P_k(G)$ if this point satisfies with equality the partition constraint (2a) (respectively (3a)) associated with $M$ when $k$ is odd (respectively even). Moreover, $M$ is tight for a face $F$ of $P_k(G)$ if it is tight for all the points of $F$.

The following results give some insights on the structure of tight multicuts.

Theorem 2.10 [17, Theorem 2.3 and Lemma 3.1] Let $x$ be a point of $P_{2h+1}(G)$, and let $M = \delta(V_1, \ldots, V_{dM})$ be a multicut tight for $x$. Then, the following hold:

(i) If $d_M \geq 3$, then $x(\delta(V_i) \cap \delta(V_j)) \leq h + 1$ for all $i \neq j \in \{1, \ldots, d_M\}$.

(ii) $G \setminus V_i$ is connected for all $i = 1, \ldots, d_M$.

Lemma 2.11 Let $v$ be a degree 2 vertex of $G$ and $M$ be a multicut of $G$ strictly containing $\delta(v) = \{uv, vw\}$. If $M$ is tight for a point of $P_k(G)$ with $k \geq 2$, then both $M \setminus uv$ and $M \setminus vw$ are multicuts of $G$ of order $d_M - 1$.

Proof It suffices to show that $u$ and $w$ belong to different classes of the multicut $M = \delta(v, V_2, \ldots, V_{dM})$. Suppose that $u, w \in V_2$. Then $M$ is the union of the two multicuts $\delta(v)$ and $M' = \delta(v) \cup V_2, \ldots, V_{dM}$. Since $d_{\delta(v)} + d_{M'} = d_M + 1$, the sum of the partition inequalities associated with $\delta(v)$ and $M'$ implies that the partition inequality associated with $M$ is tight for no point of $P_k(G)$ for every $k \geq 2$. □

Chopra [10] gave sufficient conditions for an inequality to be facet-defining for $P_k(G)$. The following proposition is a direct consequence of Theorems 2.4 and 2.6 of [10].

Lemma 2.12 Let $G$ be a connected graph having a $K_4$-minor. Then, there exist two disjoint nonempty subsets of edges of $G$, $E'$ and $E''$, and a rational $b$ such that

$$
x(E') + 2x(E'') \geq b,
$$

is a facet-defining inequality of $P_{2h+1}(G)$.
Chen et al. [7] provided a box-TDI system for $P_2(G)$ for series–parallel graphs.

**Theorem 2.13** [7, Theorem 1.1] The system:

\[
\begin{align*}
\frac{1}{2} x(D) & \geq 1 \text{ for all cuts } D \text{ of } G, \\
x & \geq 0
\end{align*}
\]  

is box-TDI if and only if $G$ is a series–parallel graph.

This result proves that the polyhedron $P_2(G)$ is box-TDI for all series–parallel graphs, and gives a TDI system describing this polyhedron in this case. However, Theorem 2.13 is not sufficient to state that $P_2(G)$ is a box-TDI polyhedron if and only if $G$ is series–parallel.

### 3 Box-TDIness of $P_k(G)$

In this section we show that, for $k \geq 2$, $P_k(G)$ is a box-TDI polyhedron for a connected graph $G$ if and only if $G$ is series–parallel. Since $P_k(G) = \emptyset$ when $G$ is not connected, we assume from now on that $G$ is connected.

When $k \geq 2$, $P_k(G)$ is not always box-TDI, as stated in Lemma 3.1. Indeed, by Theorem 2.4, if a polyhedron has a nonequimodular face-defining matrix, then it is not box-TDI. The proof of Lemma 3.1 exhibits such a matrix when $G$ has a $K_4$-minor. This follows from the existence of a particular facet-defining inequality when $k$ is odd, as shown by Chopra [10]. When $k$ is even, we build a nonequimodular face-defining matrix based on the structure of cuts in a $K_4$-minor.

**Lemma 3.1** For $k \geq 2$, if $G = (V, E)$ has a $K_4$-minor, then $P_k(G)$ is not box-TDI.

**Proof** When $k = 2h + 1$ is odd, Lemma 2.12 shows that there exists a facet-defining inequality that is described by a nonequimodular matrix as $P_k(G)$ is full-dimensional. Thus, $P_k(G)$ is not box-TDI by Statement (ii) of Theorem 2.4.

We now prove the case when $k$ is even. Since $G$ has a $K_4$-minor, there exists a partition $\{V_1, \ldots, V_4\}$ of $V$ such that $G[V_i]$ is connected and $\delta(V_i, V_j) \neq \emptyset$ for all $i < j \in \{1, \ldots, 4\}$. We now prove that the matrix $A$ whose three rows are $\chi_{\delta(V_i)}$ for $i = 1, 2, 3$ is a face-defining matrix of $P_k(G)$ which is not equimodular. This will end the proof by Statement (ii) of Theorem 2.4.

Let $e_{ij}$ be an edge in $\delta(V_i, V_j)$ for all $i < j \in \{1, \ldots, 4\}$. The submatrix of $A$ formed by the columns associated with edges $e_{ij}$ is the following:

\[
\begin{bmatrix}
1 & e_{12} & e_{13} & e_{14} & e_{24} & e_{34} \\
1 & 1 & 0 & 1 & 0 & 0 \\
1 & 0 & 1 & 0 & 1 & 0 \\
0 & 1 & 1 & 0 & 1 & 0
\end{bmatrix}
\]

The matrix $A$ is not equimodular as the first three columns form a matrix of determinant -2 whereas the last three ones give a matrix of determinant 1.
By Observation 2.7, to show that A is face-defining, it is enough to exhibit \(|E| - 2\) affinely independent points of \(P_k(G)\) satisfying \(x(\delta(V_i)) = k\) for \(i = 1, 2, 3\).

Let \(D_1 = \{e_{12}, e_{14}, e_{23}, e_{34}\}\), \(D_2 = \{e_{12}, e_{13}, e_{24}, e_{34}\}\), \(D_3 = \{e_{13}, e_{14}, e_{23}, e_{24}\}\) and \(D_4 = \{e_{14}, e_{24}, e_{34}\}\). First, we define the points \(S_j = \sum_{i=1}^{4} k \chi(V_i) + \frac{k}{2} \chi(D_j)\), for \(j = 1, 2, 3\), and \(S_4 = \sum_{i=1}^{4} k \chi(V_i) + k \chi(D_4)\). Note that they are affinely independent.

Now, for each edge \(e \notin \{e_{12}, e_{13}, e_{14}, e_{23}, e_{24}, e_{34}\}\), we construct the point \(S_e\) as follows. When \(e \in \delta(V_i)\) for some \(i = 1, \ldots, 4\), we define \(S_e = S_4 + \chi^e\). Adding the point \(S_e\) maintains affine independence as \(S_e\) is the only point not satisfying \(x_e = k\).

When \(e \in \delta(V_i, V_j)\) for some \(i, j\), we define \(S_e = S_4 + \chi^{e_{ij}} + \chi^e\), where \(S_e\) is \(S_4\) if \(e \in \delta(V_1, V_4) \cup \delta(V_2, V_3)\) and \(S_2\) otherwise. Affine independence comes because \(S_e\) is the only point involving \(e\).

In total, we built \(4 + |E| - 6 = |E| - 2\) affinely independent points. \(\square\)

The following theorem characterizes the class of graphs for which \(P_k(G)\) is box-TDI. The case \(k\) even is obtained using the box-TDIness for \(k = 2\) and the fact that integer dilations maintain box-TDIness. For the case \(k\) odd, on the contrary to what is generally done, the proof does not exhibit a box-TDI system describing \(P_k(G)\). For this case, the proof is by induction on the number of edges of \(G\). We prove that series-parallel operations preserve the box-TDIness of the polyhedron. The most technical part of the proof is the subdivision of an edge \(uw\) into two edges \(uv\) and \(vw\). We proceed by contradiction: by Theorem 2.4, we suppose that there exists a face \(F\) of \(P_k(G)\) defined by a nonequimodular matrix. We study the structure of the inequalities corresponding to this matrix. In particular, we show that they are all associated with multicut, and that these multicut contain either both \(uv\) and \(vw\), or none of them—see Claims 3.1, 3.2, and 3.3. These last results allow us to build a nonequimodular face-defining matrix for the smaller graph, which contradicts the induction hypothesis.

**Theorem 3.2** For \(k \geq 2\), \(P_k(G)\) is a box-TDI polyhedron if and only if \(G\) is series-parallel.

**Proof** Necessity follows from Lemma 3.1. Let us now prove sufficiency. When \(k = 2\), the box-TDIness of System (5) has been shown by Chen et al. [7]. This implies box-TDIness for all even \(k\): multiplying the right-hand side of a box-TDI system by a positive rational preserves its box-TDIness [34, Section 22.5]. The system obtained by multiplying by \(\frac{k}{2}\) the right-hand side of System (5) describes \(P_k(G)\) when \(k\) is even. Hence, the latter is a box-TDI polyhedron.

The rest of the proof is devoted to the case where \(k = 2h + 1\) for some \(h \geq 1\). To this end, we prove that for every face of \(P_{2h+1}(G)\) there exists an equimodular face-defining matrix. The characterization of box-TDIness given in Theorem 2.4 concludes. We proceed by induction on the number of edges of \(G\).

If \(G\) is trivial, then \(P_{2h+1}(G) = \{x \in \mathbb{R}_+^e : x \geq 2h + 1\}\) is box-TDI. If \(G\) is the circuit \(\{e, f\}\), then \(P_{2h+1}(G) = \{xe, xf \in \mathbb{R}_+^e : xe + xf \geq 2h + 1\}\) is also box-TDI.

(1-sum) Let \(G\) be the 1-sum of two series-parallel graphs \(G' = (W', E')\) and \(G'' = (W'', E'')\). By induction, there exist two box-TDI systems \(A' y \geq b'\) and \(A'' z \geq b''\) describing respectively \(P_{2h+1}(G')\) and \(P_{2h+1}(G'')\). If \(v\) is the vertex of \(G\) obtained by the identification, \(G\backslash v\) is not connected, hence, by Statement (ii) of Theorem 2.10,
a multicut $M$ of $G$ is tight for a face of $P_{2h+1}(G)$ only if $M \subseteq E'$ or $M \subseteq E''$. It follows that for every face $F$ of $P_{2h+1}(G)$ there exist faces $F'$ and $F''$ of $P_{2h+1}(G')$ and $P_{2h+1}(G'')$ respectively, such that $F = F' \times F''$. Then $P_{2h+1}(G) = \{(y, z) \in \mathbb{R}^E_+ \times \mathbb{R}^{E''}_+ : A'y \geq b', A''z \geq b''\}$ and so it is box-TDI.

(Parallelization) Let $G = (V, E)$ be obtained from a series–parallel graph $G'$ by adding an edge $g$ parallel to an edge $f$ of $G'$ and suppose that $P_{2h+1}(G')$ is box-TDI. Let $A'x \geq b'$ be a box-TDI system describing $P_{2h+1}(G')$. Note that $P_{2h+1}(G)$ is described by $Ax \geq b, x_f \geq 0, x_g \geq 0$, where $A$ is the matrix obtained by duplicating $f$’s column. By Theorem 22.10 of [34], the system $Ax \geq b$ is box-TDI, hence so is $Ax \geq b, x_f \geq 0, x_g \geq 0$. Thus, $P_{2h+1}(G)$ is a box-TDI polyhedron.

(Subdivision) Let $G = (V, E)$ be obtained by subdividing an edge $uw$ of a series–parallel graph $G' = (V', E')$ into a path of length two $uv, vw$. By contradiction, suppose there exists a nonempty face $F = \{x \in P_{2h+1}(G) : A_Fx = b_F\}$ such that $A_F$ is a face-defining matrix for $F$ which is not equimodular. Take such a face with maximum dimension. Then, every submatrix of $A_F$ which is face-defining for a face of $P_{2h+1}(G)$ is equimodular. We may assume that $A_F$ is defined by the partition constraints (2a) associated with the set of multicuts $M_F$ and the nonnegativity constraints associated with the set of edges $E_F$.

Claim 3.1 $E_F = \emptyset$.

Proof Suppose there exists an edge $e \in E_F$. Let $H = G \setminus e$ and let $A_{F_H}x = b_{F_H}$ be the system obtained from $A_Fx = b_F$ by removing the column and the nonnegativity constraint associated with $e$. Since the matrix $A_F$ is of full row rank, so is $A_{F_H}$. Since $e \in E_F$, for all multicuts $M$ tight for $F$ not containing $e$, $M \cup e$ is not a multicut. Hence $M \setminus e$ is a multicut of $H$ of order $d_M$, for all $M$ in $M_F$. Hence, the set $F_{H} = \{x \in P_{2h+1}(H) : A_{F_H}x = b_{F_H}\}$ is a face of $P_{2h+1}(H)$. Moreover, deleting $e$’s coordinate of $\text{aff}(F)$ gives $\text{aff}(F_H)$ so $A_{F_H}$ is face-defining for $F_H$. By the induction hypothesis, $A_{F_H}$ is equimodular. Since maximal invertible square submatrices of $A_F$ are in bijection with those of $A_{F_H}$ and have the same determinant in absolute value, $A_F$ is equimodular, a contradiction. 

Claim 3.2 For $e \in \{uv, vw\}$, at least one multicut of $M_F$ different from $\delta(v)$ contains $e$.

Proof By contradiction, suppose for instance that $uv$ belongs to no multicut of $M_F$ different from $\delta(v)$.

First, suppose that $\delta(v)$ does not belong to $M_F$. Then, the column of $A_F$ associated with $uv$ is zero. Let $A'_F$ be the matrix obtained from $A_F$ by removing this column. Every multicut of $G$ not containing $uv$ is a multicut of $G'$ (relabelling $vw$ by $uw$), so the rows of $A'_F$ are associated with multicut of $G'$. Thus, $F' = \{x \in P_{k}(G') : A'_Fx = b_F\}$ is a face of $P_{2h+1}(G')$. Removing $uv$’s coordinate from the points of $F$ gives a set of points of $F'$ of affine dimension at least $\dim(F) - 1$. Since $A'_F$ has the same rank as $A_F$ and has one column fewer than $A_F$, then $A'_F$ is face-defining for $F'$ by Observation 2.7. By the induction hypothesis, $A'_F$ is equimodular. Since adding a column of zeros preserves equimodularity, $A_F$ is also equimodular.

Suppose now that $\delta(v)$ belongs to $M_F$. Then, the column of $A_F$ associated with $uv$ has zeros in each row but $\chi^\delta(v)$. Let $A_F^*x = b_F^*$ be the system obtained from $A_Fx = b_F$
by removing the equation associated with $\delta(v)$. Then $F^* = \{x \in P_k(G) : A_F^*x = b_F^*\}$ is a face of $P_k(G)$ of dimension $\dim(F) + 1$. Indeed, it contains $F$ and $z + \alpha \chi_{uv} \notin F$ for every point $z$ of $F$ and $\alpha > 0$. Hence, $A_F^*$ is face-defining for $F^*$. This matrix is equimodular by the maximality assumption on $F$, and so is $A_F$ by Statement (ii) of Observation 2.6.

\begin{claim}
\label{claim:multicut-strict}
$|M \cap \delta(v)| \neq 1$ for every multicut $M \in \mathcal{M}_F$.
\end{claim}

\begin{proof}
Suppose there exists a multicut $M$ tight for $F$ such that $|M \cap \delta(v)| = 1$. Without loss of generality, suppose that $M$ contains $uv$ but not $vw$. Then, $F \subseteq \{x \in P_{2h+1}(G) : x_{vw} \geq x_{uv}\}$ because of the partition inequality (2a) associated with the multicut $M \Delta \delta(v)$. Moreover, the partition inequality associated with $\delta(v)$ and the integrality of $P_{2h+1}(G)$ imply $F \subseteq \{x \in P_{2h+1}(G) : x_{vw} \geq h + 1\}$. The proof is divided into two cases.

\begin{case}
If $F \subseteq \{x \in P_{2h+1}(G) : x_{vw} = h + 1\}$, we prove this case by exhibiting an equimodular face-defining matrix for $F$. By Observation 2.5, this implies that $A_F$ is equimodular, which contradicts the assumption on $F$.

Equality $x_{uv} = h + 1$ can be expressed as a linear combination of equations of $A_Fx = b_F$. Let $A'_Fx = b'_F$ denote the system obtained by replacing an equation of $A_Fx = b_F$ by $x_{uv} = h + 1$ in such a way that the underlying affine space remains unchanged. Denote by $\mathcal{N}$ the set of multicuts of $\mathcal{M}_F$ containing $vw$ but not $uv$. If $\mathcal{N} \neq \emptyset$, then let $N$ be in $\mathcal{N}$. We now modify the system $A'_Fx = b'_F$ by performing the following operations.

1. For all $M \in \mathcal{M}_F$ strictly containing $\delta(v)$, replace the equation associated with $M$ by the partition constraint (2a) associated with $M \setminus vw$ set to equality, that is, $x(M \setminus vw) = (h + 1)d_M \setminus vw - 1$.
2. If $\delta(v) \in \mathcal{M}_F$, then replace the equation associated with $\delta(v)$ by the constraint $x_{uv} = h$.
3. If $\mathcal{N} \neq \emptyset$, then replace the equation associated with $N$ by the constraint $x_{uv} = h + 1$.
4. For all $M \in \mathcal{N} \setminus N$, replace the equation associated with $M$ by the partition constraint (2a) associated with $M \Delta \delta(v)$ set to equality, that is, $x(M \Delta \delta(v)) = (h + 1)d_M \setminus \delta(v) - 1$.

These operations do not change the underlying affine space. Indeed, for every multicut $M$ strictly containing $\delta(v)$ and tight for $F$, the set $M \setminus vw$ is a multicut tight for $F$ by Lemma 2.11 and $F \subseteq \{x \in P_{2h+1}(G) : x_{vw} = h + 1\}$. If $\delta(v)$ is tight for $F$, then $F \subseteq \{x \in P_{2h+1}(G) : x_{uv} = h\}$ because $F \subseteq \{x \in P_{2h+1}(G) : x_{vw} = h + 1\}$. For $M \in \mathcal{N}$, by Observation 2.2, the set $M \Delta \delta(v)$ is a multicut of order $d_M$. The tightness of the constraint (2a) associated with $\delta(v)$ and the constraint (2a) associated with $M \Delta \delta(v)$ imply that $F \subseteq \{x \in P_{2h+1}(G) : x_{vw} \leq x_{uv}\}$. Since $F \subseteq \{x \in P_{2h+1}(G) : x_{vw} \geq x_{uv}\}$, we have $F \subseteq \{x \in P_{2h+1}(G) : x_{vw} = h + 1\}$ and $M \Delta \delta(v)$ is tight for $F$. It follows that, if $\delta(v) \in \mathcal{M}_F$, then $\mathcal{N} = \emptyset$. Therefore, at most one among Operations 2 and 3 is applied so the rank of the matrix remains unchanged.

Let $A''Fx = b''_F$ be the system obtained by removing the equation $x_{vw} = h + 1$ from $A'_Fx = b'_F$. By construction, $A''Fx = b''_F$ is composed of constraints (2a) set to equality and possibly $x_{uv} = h$ or $x_{uv} = h + 1$. Moreover, the column of $A''_F$ associated
with $vw$ is zero. Let $F'' = \{ x \in P_{2h+1}(G) : A'_{F,x} = b''_F \}$. For every point $z$ of $F$ and $\alpha \geq 0, z + \alpha \chi^{uv}$ belongs to $F''$ because the column of $A'_{F,x}$ associated with $vw$ is zero, and $z + \alpha \chi^{uv} \in P_{2h+1}(G)$. This implies that $\dim(F'') \geq \dim(F) + 1$.

If $F''$ is a face of $P_{2h+1}(G)$, then $A'_{F,x}$ is face-defining for $F''$ by Observation 2.7 and because $A'_{F}$ is face-defining for $F$. By the maximality assumption on $F$, $A'_{F,x}$ is equimodular, and hence so is $A'_{F}$ by Statement (i) of Observation 2.6.

Otherwise, by construction, $F'' = F^* \cap \{ x \in \mathbb{R}^E : x_{uv} = t \}$ where $F^*$ is a face of $P_{2h+1}(G)$ strictly containing $F$ and $t \in \{ h, h + 1 \}$. Therefore, there exists a face-defining matrix for $F''$ given by a face-defining matrix for $F^*$ and the row $\chi^{uv}$. Such a matrix is equimodular by the maximality assumption of $F$ and Statement (i) of Observation 2.6. Hence, $A''_{F}$ is equimodular by Observation 2.5, and so is $A'_{F}$ by Statement (i) of Observation 2.6.

Case 2 $F \not
\subseteq \{ x \in P_{2h+1}(G) : x_{uv} = h + 1 \}$. Thus, there exists $z \in F$ such that $z_{uv} > h + 1$. By Claim 3.2, there exists a multicut $N \neq \delta(v)$ containing $vw$ which is tight for $F$. By Statement (i) of Theorem 2.10, the existence of $z$ implies that $N$ is a bond, hence it does not contain $uv$. The set $L = N \Delta \delta(v)$ is a bond of $G$. The partition inequality (2a) associated with $L$ implies that $F \subseteq \{ x \in P_{2h+1}(G) : x_{uv} = x_{uv} \}$ and $L$ is tight for $F$. Moreover, $N$ is the unique multicut tight for $F$ containing $vw$.

Suppose indeed that there exists a multicut $B$ containing $vw$ tight for $F$. Then, $B$ is a bond by Statement (i) of Theorem 2.10 and the existence of $z$. Moreover, $B \Delta N$ is a multicut not containing $vw$. This implies that no point of $F$ satisfies the partition constraint associated with $B \Delta N$ because $x(B \Delta N) = x(B) + x(N) - 2x(B \cap N) = 2(2h + 1) - 2x(B \cap N) \leq 4h + 2 - 2x_{uv} \leq 2h$, a contradiction.

Consider the matrix $A'_{F}$ obtained from $A_{F}$ by removing the row associated with $N$. Matrix $A'_{F}$ is a face-defining matrix for a face $F^* \supseteq F$ of $P_{2h+1}(G)$ because $F^*$ contains $F$ and $z + \alpha \chi^{uv}$ for every point $z$ of $F$ and $\alpha > 0$. By the maximality assumption, the matrix $A'_{F}$ is equimodular. Let $B_{F}$ be the matrix obtained from $A_{F}$ by replacing the row $\chi^{N}$ by the row $\chi^{N} - \chi^{L}$. Then, $B_{F}$ is face-defining for $F$. Moreover, $B_{F}$ is equimodular by Statement (ii) of Observation 2.6—a contradiction. □

Let $A'_{F,x} = b'_{F}$ be the system obtained from $A_{F,x} = b_{F}$ by removing $uv$’s column from $A_{F}$ and subtracting $h + 1$ times this column to $b_{F}$. We now show that $\{ x \in P_{2h+1}(G') : A'_{F,x} = b'_{F} \}$ is a face of $P_{2h+1}(G')$ if $\delta(v) \notin M_{F}$, and of $P_{2h+1}(G') \cap \{ x : x_{u \bar{u} v} = h \}$ otherwise. Indeed, consider a multicut $M$ in $M_{F}$. If $M = \delta(v)$, then the equation of $A'_{F,x} = b'_{F}$ induced by $M$ is nothing but $x_{u \bar{u} v} = h$. Otherwise, by Lemma 2.11 and Claim 3.3, the set $M \setminus uv$ is a multicut of $G'$ (relabeling $vw$ by $uv$) of order $d_{M}$ if $uv \notin M$ and $d_{M} - 1$ otherwise. Thus, the equation of $A'_{F,x} = b'_{F}$ induced by $M$ is the partition constraint (2a) associated with $M \setminus uv$ set to equality.

By construction and Claim 3.3, $A'_{F}$ has full row rank and one column less than $A_{F}$. We prove that $A'_{F}$ is face-defining by exhibiting $\dim(F)$ affinely independent points of $P_{2h+1}(G')$ satisfying $A'_{F,x} = b'_{F}$. Because of the integrality of $P_{2h+1}(G)$, there exist $n = \dim(F) + 1$ affinely independent integer points $z^1, \ldots, z^n$ of $F$. By Claims 3.2 and 3.3, there exists a multicut strictly containing $\delta(v)$. Then, Statement (i) of Theorem 2.10 implies that $F \subseteq \{ x \in \mathbb{R}^E : x_{uv} \leq h + 1, x_{u \bar{u} v} \leq h + 1 \}$. Combined with the partition inequality $x_{u \bar{u} v} + x_{\bar{u} v} \geq 2h + 1$ associated with $\delta(v)$, this implies that at least one of $\sum_{i=1}^{n} z_{u \bar{u} v}^{i}$, and $z_{u \bar{u} v}^{i}$ is equal to $h + 1$ for $i = 1, \ldots, n$. Since exchanging the $uv$
and $v w$ coordinates of any point of $F$ gives a point of $F$ by Claim 3.3, the hypotheses on $z^1, \ldots, z^n$ are preserved under the assumption that $z^i_{uv} = h + 1$ for $i = 1, \ldots, n - 1$. Let $y^1, \ldots, y^{n-1}$ be the points obtained from $z^1, \ldots, z^{n-1}$ by removing $uv$'s coordinate.

Since every multicut of $G'$ is a multicut of $G$ with the same order, $y^1, \ldots, y^{n-1}$ belong to $P_{2h+1}(G')$. By construction, they satisfy $A'Fy = b'$ so they belong to a face of $P_{2h+1}(G')$ or $P_{2h+1}(G') \cap \{x : x_{uw} = h\}$. This implies that $A'F$ is a face-defining matrix of $P_{2h+1}(G')$ if $\delta(v) / \in M$, and of $P_{2h+1}(G') \cap \{x : x_{uw} = h\}$ otherwise.

By induction, $P_{2h+1}(G')$ is a box-TDI polyhedron and hence so is $P_{2h+1}(G') \cap \{x : x_{uw} = h\}$. Hence, $A'F$ is equimodular by Theorem 2.4. Since $AF$ is obtained from $A'F$ by copying a column, then $AF$ is also equimodular—a contradiction. $\square$

By definition of box-TDIness and $Q_k(G)$, Theorem 3.2 implies that $Q_k(G)$ is box-TDI when $G$ is series–parallel. The converse does not hold. Indeed, for instance, when $G = (V, E)$ is a minimal $k$-edge-connected graph, $Q_k(G)$ is nothing but the single point $\chi^E$ so it is a box-TDI polyhedron.

4 An integer TDI system for $P_{2h}(G)$

Let $G$ be a series–parallel graph. In this section we provide an integer TDI system for $P_{2h}(G)$ with $h$ positive and integer.

The proof of the main result of this section is based on the characterization of TDIness by means of Hilbert bases. A set of vectors $\{v^1, \ldots, v^k\}$ is a Hilbert basis if each integer vector that is a nonnegative combination of $v^1, \ldots, v^k$ can be expressed as a nonnegative integer combination of them. The link between Hilbert basis and TDIness is stated in the following theorem.

**Theorem 4.1** (Theorem 22.5 of [34]) A system $Ax \geq b$ is TDI if and only if for every face $F$ of $P = \{x : Ax \geq b\}$, the rows of $A$ associated with tight constraints for $F$ form a Hilbert basis.

In the previous theorem, we could restrict to minimal faces: indeed, the cone generated by the constraints tight for a face $F$ is a face of the cone generated by the constraints active for a face $F' \subseteq F$ [34].

**Remark 4.2** A system $Ax \geq b$ is TDI if and only if, for each minimal face $F$ of $P = \{x : Ax \geq b\}$, the rows of $A$ associated with constraints tight for $F$ form a Hilbert basis.

The rest of the section is devoted to prove that the system given by the partition constraints and the nonnegativity constraints, which describes $P_k(G)$ when $k$ is even, is TDI when $G$ is series–parallel.

The proof is based on the TDIness of System (5) and the structure of inequalities (3a). Their right-hand sides are proportional to $h$, hence it is enough to prove the case $k = 2$. This allows us to use Theorem 2.13 to obtain a TDI system for $P_2(G)$. In terms of Hilbert bases, the TDIness of this system implies that, given a face $F$ of $P_2(G)$, the integer points of the associated cone are the half sum of the cuts tight for
Theorem 4.3  For a series–parallel graph $G$ and a positive integer $h$, System (3) is TDI.

Proof  We only prove the case $h = 1$ since multiplying the right hand side of a system by a positive constant preserves its TDIness [34, Section 22.5].

The proof is done by induction on the number of edges of the graph $G = (V, E)$. When $G$ consists of two vertices connected by a single edge $e$, System (3) is $x_e \geq 2, x_e \geq 0$ and is TDI. If $G$ is the circuit $\{e, f\}$, System (3) is $x_e + x_f \geq 2, x \geq 0$ and is TDI.

(Parallelization) Let now $G$ be obtained from a series–parallel graph $H$ by adding an edge $g$ parallel to an edge $f$ of $H$. System (3) associated with $G$ is obtained from that associated with $H$ by duplicating $f$'s column in constraints (3a) and adding the nonnegativity constraint $x_g \geq 0$. By Lemma 3.1 of [7], System (3) is TDI.

For the other cases, we prove the TDIness of System (3) associated with $G$ using Remark 4.2. More precisely, we prove that for any extreme point $z$ of $P_2(G)$, the set of vectors $\{x^M : M \in T_z\} \cup \{x^e : e \in E, z_e = 0\}$ is a Hilbert basis, where $T_z$ is the set of multicuts tight for $z$.

(1-sum) Let $G$ be the 1-sum of two series–parallel graphs $G^1 = (W^1, E^1)$ and $G^2 = (W^2, E^2)$ and let $z$ be an extreme point of $P_2(G)$. By construction, we have $z = (z_1^1, z_2^2)$ where $z^i = P_2(G^i)$ for $i = 1, 2$. Moreover, for each multicut $M \in T_z$, the graph obtained from $G(z)$ by contracting the edges of $E \setminus M$ is a circuit. Indeed, it is 2-edge-connected since $G(z)$ is, and it has $z(M) = d_M$ edges and $d_M$ vertices. Therefore $M$ is either a multicut of $G^1$ tight for $z^1$ or one of $G^2$ tight for $z^2$.

By induction, Systems (3) associated with $G^1$ and $G^2$ are TDI. Thus, $\{x^M : M \in T_z \cap M(G^i)\} \cup \{x^e : e \in E^i, z_e = 0\}$ is a Hilbert basis for $i = 1, 2$ by Theorem 4.1. Since they belong to disjoint spaces, their union is a Hilbert basis. By Theorem 4.1, System (3) is TDI.

(Subdivision) Let $G = (V, E)$ be obtained by subdividing an edge $uv$ of a series–parallel graph $G' = (V', E')$ into a path of length two $uv, vw$, and let $z$ be an extreme point of $P_2(G)$.

Without loss of generality, suppose $z_{uv} \geq z_{vw}$. Define $z' \in ZE'$ by $z_{uv}' = z_{vw}$ and $z_e' = z_e$ for all edges $e$ in $E' \setminus uv$. Note that $z'$ belongs to $P_2(G')$ since $G'(z')$ is obtained by contracting the edge $uv$ in $G(z)$, and this contraction preserves 2-edge-connectivity.

Note that for all $e \in E, z_e \in \{0, 1, 2\}$. Indeed, since $z$ is an extreme point of $P_2(G)$ which is also described by System (1), if $z_e > 0$, then $e$ belongs to a cut $D$ tight for $z$. Moreover, as $z_{uv} \geq z_{vw}$, the partition constraint (3a) associated with $\delta(u)$ implies that $z_{uv} \in \{1, 2\}$. We now consider two different cases depending on the value of $z_{uv}$.

Case 1 $z_{uv} = 2$. We first show that every multicut of $T_z$ containing $uv$ is a bond. Indeed, note that every multicut $M$ with $d_M = 2$ is a bond. If a multicut $M = \delta(V_1, \ldots, V_{d_M}) \in T_z$ satisfies $d_M \geq 3$ and $uv \in \delta(V_1, V_2)$, then $M' = \delta(V_1 \cup V_2, V_3, \ldots, V_{d_M})$ is a bond.
multicut and satisfies
\[ z(M') \leq z(M) - 2 < d_M - 1 = d_{M'} . \]

Hence, the partition constraint (3a) associated with \( M' \) is violated, a contradiction.

Moreover, there exists at most one bond of \( T_z \), say \( N \), containing \( uv \). As otherwise suppose there exist two bonds \( B_1 \) and \( B_2 \) in \( T_z \) containing \( uv \). Then, \( z(B_1 \cup B_2) \leq z(B_1) + z(B_2) - 2z_{uv} = 0 \), which contradicts the constraint (3a) associated with the multicut \( B_1 \cup B_2 \). For a multicut \( M \) not containing \( uv \), \( M \in T_z \) if and only if \( M \in T_{z'} \). This implies that \( T_z = T_{z'} \cup N \). By induction and Theorem 4.1, \( T_{z'} \cup E_{z'} \) is a Hilbert basis. As \( E_z = E_{z'} \) (identifying \( uv \) and \( vw \)) and \( N \) is the only member of \( T_z \cup E_z \) containing \( uv \), \( T_z \cup E_z \) is also a Hilbert basis.

Case 2 \( z_{uv} = 1 \). Let \( v \) be an integer point of the cone generated by \( T_z \cup E_z \). We prove that \( v \) can be expressed as an integer nonnegative combination of the vectors of \( T_z \cup E_z \). This implies that \( T_z \cup E_z \) is a Hilbert basis.

Let \( B_z \) be the set of bonds of \( T_z \). Since System (5) is a TDI system describing \( P_2(G) \) in series–parallel graphs, the set of vectors \( \{ \chi_B : B \in B_z \cup E_z \} \) forms a Hilbert basis by Theorem 4.1. Then, there exist \( \lambda_B \in \frac{1}{2}\mathbb{Z}_+ \) for all \( B \in B_z \) and \( \mu_e \in \mathbb{Z}_+ \) for all \( e \in E_z \) such that \( v = \sum_{B \in B_z} \lambda_B \chi_B + \sum_{e \in E_z} \mu_e \chi_e \).

Since \( z_{uv} \geq z_{vw} \), the partition inequality (3a) associated with \( \delta(v) \) implies that \( z_{uw} = 1 \) and \( \delta(v) \in T_z \). In particular, \( vw \notin E_{z'}. \) The vector \( v \) is an integer combination of vectors of \( T_z \cup E_z \) if and only if \( v - [\lambda_{\delta(v)}] \chi_{\delta(v)} \) is, thus we may assume that \( \lambda_{\delta(v)} \in \{0, \frac{1}{2}\} \). Define \( w \in \mathbb{Z}^{E'} \) by:

\[
w_e = \begin{cases} v_{uv} + v_{vw} - 2\lambda_{\delta(v)} & \text{if } e = uw, \\ v_e & \text{otherwise.} \end{cases}
\]

Note that \( (B \setminus uw) \cup uv \) and \( (B \setminus uw) \cup vw \) are bonds of \( T_z \) whenever \( B \) is a bond of \( T_{z'} \) containing \( uv \) because \( z'_{uw} = z_{uw} = z_{vw} = 1 \). Moreover, a bond \( B \) of \( T_{z'} \) which does not contain \( uw \) is a bond of \( T_z \). Since \( \delta(v) \) is the unique bond of \( G \) containing both \( uv \) and \( vw \) and \( E_z = E_{z'} \), we have:

\[
w = \sum_{B \in B_z : uw \in B} (\lambda_{(B \setminus uw) \cup uv} + \lambda_{(B \setminus uw) \cup vw}) \chi_B + \sum_{B \in B_{z'} : uw \notin B} \lambda_B \chi_B + \sum_{e \in E_z} \mu_e \chi_e .
\]

Thus, \( w \) belongs to the cone generated by \( T_{z'} \cup E_{z'} \). By the induction hypothesis, \( T_{z'} \cup E_{z'} \) is a Hilbert basis, hence there exist \( \lambda'_{M} \in \mathbb{Z}_+ \) for all \( M \in T_{z'} \) and \( \mu'_e \in \mathbb{Z}_+ \) for all \( e \in E_{z'} \) such that \( w = \sum_{M \in T_{z'}} \lambda'_M \chi_M + \sum_{e \in E_{z'}} \mu'_e \chi_e \).

Consider the family \( \mathcal{N} \) of multicut of \( T_{z'} \) where each multicut \( M \) of \( T_{z'} \) appears \( \lambda'_{M} \) times. Suppose first that \( \lambda_{\delta(v)} = 0 \). Then, \( v_{uv} + v_{vw} \) multicut of \( \mathcal{N} \) contain \( uw \). Let \( \mathcal{P} \) be a family of \( v_{uv} \) multicut of \( \mathcal{N} \) containing \( uv \) and \( Q = \{ F \in \mathcal{N} : uw \in F \} \setminus \mathcal{P} \). Then, we have:

\[
v = \sum_{M \in \mathcal{N} : uw \notin M} \chi_M + \sum_{M \in \mathcal{P}} \chi_{(M \setminus uw) \cup uv} + \sum_{M \in Q} \chi_{(M \setminus uw) \cup vw} + \sum_{e \in E_{z'}} \mu'_e \chi_e . \quad (6)
\]
Suppose now that $\lambda_\delta(v) = \frac{1}{2}$. Then, $w_{uv} = v_{uv} + v_{vw} - 1$ multicuts of $N$ contain $uw$. Let $P$ be a family of $v_{uv} - 1$ multicuts of $N$ containing $uw$, let $Q$ be a family of $v_{vw} - 1$ multicuts in $\{ F \in N : uw \in F \} \setminus P$, and denote by $N$ the unique multicut of $N$ containing $uw$ which is not in $P \cup Q$. Then, we have

$$v = \sum_{M \in N : uw \notin M} \chi^M + \sum_{M \in P} \chi^{(M \setminus uw) \cup uv} + \sum_{M \in Q} \chi^{(M \setminus uw) \cup vw} + \chi^{(N \setminus uw) \cup \delta(v)}$$

(7)

Every $M \in T_z'$ not containing $uw$ is in $T_z$. For every $M \in T_z'$ containing $uw$, $(M \setminus uw) \cup uv$, $(M \setminus uw) \cup vw$ and $(M \setminus uw) \cup \delta(v)$ belong to $T_z$ since $z'_{uw} = z_{uv} = z_{vw} = 1$. Since $\mathcal{E}_z = \mathcal{E}_z'$, then $v$ is a nonnegative integer combination of vectors of $T_z \cup \mathcal{E}_z$ in both (6) and (7). This proves that $T_z \cup \mathcal{E}_z$ is a Hilbert basis. Therefore by Remark 4.2, System (3) is TDI.

The box-TDIness of $P_k(G)$ and the TDIness of System (3) give the following result.

**Corollary 4.4** System (3) is box-TDI if and only if $G$ is series–parallel.

**Proof** If $G$ is series–parallel, then System (3) is box-TDI by Theorems 3.2 and 4.3, since a TDI system describing a box-TDI polyhedron is box-TDI [12]. If $G$ is not series–parallel, Theorem 3.2 ensures that $P_k(G)$ is not box-TDI, therefore System (3) is not box-TDI.

Theorem 4.3 leaves open the following problem:

**Open Problem 4.5** Characterize the classes of graphs such that System (3) is TDI.

### 5 An integer TDI system for $P_{2h+1}(G)$

In this section, we prove that System (2) is TDI if and only if $G$ is a series–parallel graph—see Theorem 5.1. Proving the TDIness for $k$ odd is considerably more involved than for $k$ even. The first difference with the even case is the lack of a known TDI system describing $P_k(G)$ when $k$ is odd, even a noninteger one. Thus, no property of the Hilbert bases associated with $P_k(G)$ is known, and the approach used to prove Theorem 4.3 cannot be applied. Instead, following the definition of TDIness, we prove the existence of an integer optimal solution to each feasible dual problem.

Another difference with the case $k$ even follows from the structure of the partition inequalities (2a). In particular, the presence of the constant “$-1$” in the right-hand sides perturbs the structure of tight multicuts. Indeed, when $k$ is odd, the tightness of $\delta(V_1, \ldots, V_n)$ does not imply that of $\delta(V_1), \ldots, \delta(V_n)$. Consequently, it is not clear how the contraction of an edge impacts the tightness of a multicut $\delta(V_1, \ldots, V_n)$: merging adjacent $V_i$’s is not sufficient to obtain new tight multicuts. Due to the link between tight multicuts and positive dual variables, the structure of the optimal solutions to the dual problem is completely modified when subdividing an edge. Proving directly that
subdivision preserves TDI\textsuperscript{ness} turned out to be challenging, and we overcome this difficulty by deriving new properties of series–parallel graphs—see Lemma 2.3. The proof of Theorem 5.1 focuses on properties of vertices of degree 2 in a minimal counterexample to the TDI\textsuperscript{ness} of System (2). In particular, we prove that no two vertices of degree 2 are adjacent (Claim 5.9), or in the same circuit of length 4 (Claim 5.11). Moreover, no triangle contains vertices of degree 2 (Claim 5.10). By Lemma 2.3, this implies that the graph is not series–parallel. To derive these properties, we study the interplay between cuts associated with degree 2 vertices and dual optimal solutions—see Claims 5.3–5.8.

**Theorem 5.1** For $h$ positive and integer, System (2) is TDI if and only if $G$ is series–parallel.

**Proof** If $G$ is not series–parallel, then System (2) is not TDI because every TDI system with integer right-hand side describes an integer polyhedron \cite{22}, but when $G$ has a $K_4$-minor, System (2) describes a noninteger polyhedron \cite{10}.

We now prove that, if $G$ is series–parallel, then System (2) is TDI. We prove the result by contradiction. Let $G = (V, E)$ be a series–parallel graph such that System (2) is not TDI. By definition of TDI\textsuperscript{ness}, there exists $c \in \mathbb{Z}^E$ such that $\mathcal{D}(G, c)$:

\[
\max \sum_{M \in \mathcal{M}_G} b_M y_M \\
\text{s.t.} \\
\sum_{M \in \mathcal{M}_G: e \in M} y_M \leq c_e \quad \text{for all } e \in E, \quad (8a) \\
y_M \geq 0 \quad \text{for all } M \in \mathcal{M}_G, \quad (8b)
\]

is feasible, bounded, but admits no integer optimal solution, where $b_M = (h+1)d_M - 1$ for all $M \in \mathcal{M}_G$. Without loss of generality, we assume that:

(i) $G$ has a minimum number of edges,
(ii) $\sum_{e \in E} c_e$ is minimum with respect to (i).

By definition, $\mathcal{D}(G, c)$ is feasible if and only if $c \geq 0$. Hence, by minimality assumption (ii), $\mathcal{D}(G, c')$ has an optimal integer solution for every integer $c' \neq c$ such that $0 \leq c' \leq c$.

Let $M$ be a multicut of $G$. We denote by $\xi_M$ the vector of $\{0, 1\}^\mathcal{M}_G$ whose only nonzero coordinate is the one associated with $M$. We say that $M$ is active for a solution $y$ to $\mathcal{D}(G, c)$ if $y_M > 0$. Note that, by complementary slackness, a multicut is active for an optimal solution to $\mathcal{D}(G, c)$ only if it is tight for an optimal solution to the primal problem. In particular, if a multicut is tight for no point of $P_{2h+1}(G)$, then it is active for no optimal solution to $\mathcal{D}(G, c)$. Thus, we will use Lemma 2.11 and Theorem 2.10 to deduce properties on the optimal solutions to $\mathcal{D}(G, c)$.

**Claim 5.1** $G$ is simple, 2-connected, and nontrivial.

**Proof** Suppose for a contradiction that there exist two parallel edges $e_1$ and $e_2$ and $c_{e_1} \leq c_{e_2}$. Since a multicut contains either both $e_1$ and $e_2$ or none of them, the inequality

\[ c_{e_1} \leq c_{e_2} \]
admits an integer optimal solution $y$ is a solution to $D$ associated with $e$. By Statement (ii) of Theorem 2.10, the multicuts of $G$ that intersect both $E_1$ and $E_2$ are not tight for the points of $C_{2h+1}(G)$, by complementary slackness, these multicuts are not active for the optimal solutions to $D(G,c)$. Hence, every optimal solution $y$ to $D(G,c)$ is of the form:

$$y_M = \begin{cases} 
  y^1_M & \text{if } M \in \mathcal{M}_{G_1}, \\
  y^2_M & \text{if } M \in \mathcal{M}_{G_2}, \\
  0 & \text{otherwise}, 
\end{cases}$$

where $y^i$ is an optimal solution to $D((G_i,c_{|E_i}))$ for $i = 1, 2$. By minimality assumption (i), there exists an integer optimal solution $\tilde{y}^i$ to $D((G_i,c_{|E_i}))$ for $i = 1, 2$, implying that $(\tilde{y}^1, \tilde{y}^2)$ is an integer optimal solution to $D(G,c)$, a contradiction.

Finally, if $G = K_2$, $\mathcal{M}_G$ contains only one multicut, say $\{e\}$, and the optimal solution to $D(G,c)$ is $y^*_{\{e\}} = c_e$ which is integer. \qed

Claim 5.2 For all edges $e \in E$, $c_e \geq 1$.

Proof By hypothesis, $c \geq 0$ is integer and $D(G,c)$ has an optimal solution, say $y^*$. Suppose for a contradiction that there exists an edge $e \in E$ with $c_e = 0$. Set $G' = G/e$ and $c' = c_{(E\setminus e)}$. The active multicuts for $y^*$ do not contain the edge $e$ so they are multicuts of $G'$ since $\mathcal{M}_{G'} = \{M \in \mathcal{M}_G : e \notin M\}$. Hence, the point $y' \in \mathbb{R}^{\mathcal{M}_{G'}}$ defined by $y'_M = y^*_M$ for all $M \in \mathcal{M}_{G'}$ is a solution to $D(G',c')$.

By minimality assumption (i), there exists an integer optimal solution $\tilde{y}$ to $D(G',c')$. Extending $\tilde{y}$ to a point of $\mathbb{Z}^{\mathcal{M}_G}$ by setting to 0 the missing component gives an integer solution to $D(G,c)$ with cost $b^T \tilde{y} \geq b^T y^* = b^T y^*$. This is an integer optimal solution to $D(G,c)$ since $y^*$ is optimal, a contradiction to the hypothesis that $D(G,c)$ has no integer optimal solution. \qed

Claim 5.3 Every optimal solution $y$ to $D(G,c)$ satisfies $0 \leq y_M < 1$ for all $M \in \mathcal{M}_G$.

Proof By contradiction, suppose that $y^*$ is an optimal solution to $D(G,c)$ such that there exists a multicut $M$ such that $y^*_M \geq 1$. Therefore, the point $y'$ defined by $y' = y^* - \xi_M$ is a solution to $D(G,c')$ where $c' = c - \chi_M$. By minimality assumption (ii), $D(G,c')$ admits an integer optimal solution $y''$. The point $\tilde{y}$ defined by $\tilde{y} = y'' + \xi_M$ is an integer solution to $D(G,c)$ and we have:

$$b^T \tilde{y} = b^T y'' + b_M \geq b^T y' + b_M = b^T y^*.$$ 

Therefore $\tilde{y}$ is an integer optimal solution to $D(G,c)$, a contradiction. \qed

From the definition of series–parallel graphs, Claim 5.1 implies that $G$ contains at least one degree 2 vertex. Let $\tilde{V}$ be the set of degree 2 vertices of $G$. 

\[ \text{Springer} \]
Claim 5.4 Let \( v \in \widehat{V} \), \( \delta(v) = \{e_1, e_2\} \), \( y \) be an optimal solution to \( D_{(G,c)} \), and \( M_1 \) be an active multicut for \( y \) such that \( M_1 \cap \delta(v) = e_1 \). If \( \delta(v) \) is active for \( y \), then no multicut whose intersection with \( \delta(v) \) is \( e_2 \) is active for \( y \).

Proof We prove the result by contradiction. Assume that \( M_1 \) and \( \delta(v) \) are active for \( y \) and that there exists a \( M_2 \) active for \( y \) with \( M_2 \cap \delta(v) = e_2 \). By Observation 2.2, \( M_1' = M_1 \cup \delta(v) \) is a multicut of \( G \) such that \( d_{M_1'} = d_{M_i} + 1 \) for \( i = 1, 2 \). Let \( 0 < \varepsilon \leq \min\{y_{M_1}, y_{M_2}, y_{\delta(v)}\} \). Then, the point:

\[
y' = y - \varepsilon \left( \chi^{M_1} + \chi^{M_2} + \chi^{\delta(v)} \right) + \varepsilon \left( \chi^{M_1'} + \chi^{M_2'} \right)
\]

is a solution to \( D_{(G,c)} \), and we have \( b^\top y' = b^\top y + \varepsilon \), implying that \( y \) is not optimal, a contradiction. \( \square \)

Claim 5.5 For every optimal solution to \( D_{(G,c)} \), the constraints (8a) associated with the edges incident to a degree 2 vertex are tight.

Proof We prove the result by contradiction. Suppose that there exist an optimal solution \( y^* \) to \( D_{(G,c)} \) and a vertex \( v \) with \( \delta(v) = \{e_1, e_2\} \) such that the inequality (8a) associated with \( e_1 \) is not tight. For \( i = 1, 2 \), let \( s_i \) be the slack of the constraint associated with \( e_i \), that is,

\[
s_i = c_{e_i} - \sum_{M \in \mathcal{M}_G : c_{e_i} \in M} y^*_M.
\]

Inequality (8a) associated with \( e_2 \) is tight, as otherwise there exists \( 0 < \eta \leq \min\{s_1, s_2\} \), such that \( y^* + \eta \xi_{\delta(v)} \) is a solution to \( D_{(G,c)} \), a contradiction to the optimality of \( y^* \). Hence, Claims 5.2 and 5.3 imply that there are at least two distinct multcuts \( M_1 \) and \( M_2 \) active for \( y^* \) and containing \( e_2 \). Let \( 0 < \varepsilon \leq \min\{y_{M_1}, y_{M_2}, s_1\} \). For \( i = 1, 2 \), \( e_1 \in M_i \), as otherwise \( y_i = y^* + \varepsilon (\xi_{M_i \cup e_1} - \xi_{M_i}) \) is a solution to \( D_{(G,c)} \). This solution is such that \( b^\top y_i = b^\top y^* + \varepsilon (h + 1) > b^\top y^* \), a contradiction to the optimality of \( y^* \). Thus, both \( M_1 \) and \( M_2 \) contain \( \delta(v) \). Since they are distinct, at least one of them, say \( M_1 \), strictly contains \( \delta(v) \). Then, \( y'' = y^* + \varepsilon (-\xi_{M_1} + \xi_{M_1 \setminus e_2} + \xi_{\delta(v)}) \) is a solution to \( D_{(G,c)} \) because \( M_1 \setminus e_2 \) belongs to \( \mathcal{M}_G \) by Lemma 2.11. Then, \( b^\top y'' = b^\top y^* + \varepsilon (-b_{M_1} + b_{M_1} - (h + 1) + 2h + 1) > b^\top y^* \), a contradiction. \( \square \)

Given a solution \( y \) to \( D_{(G,c)} \), we define for each vertex \( v \in \widehat{V} \) the set \( \mathcal{A}_v^y \) as the set of multcuts active for \( y \) that strictly contain \( \delta(v) \). Moreover we define the value \( \alpha_v^y \) as:

\[
\alpha_v^y = \sum_{M \in \mathcal{A}_v^y} y_M.
\]

Claim 5.6 Every optimal solution \( y \) to \( D_{(G,c)} \) satisfies \( 0 < \alpha_v^y < 1 \) for all \( v \in \widehat{V} \).
Proof Suppose for a contradiction that there exist an optimal solution $y^*$ to $D_{(G,c)}$ and a vertex $v$ of $\hat{V}$ such that either $\alpha_v^y \geq 1$ or $\alpha_v^{y^*} = 0$. Denote the two edges incident to $v$ by $e_1$ and $e_2$ in such a way that $c_{e_1} \leq c_{e_2}$.

Suppose first that $\alpha_v^{y^*} \geq 1$. By Claim 5.3, there exist at least two multcuts in $A^*_v$. Let $A^*_v = \{M_1, \ldots, M_n\}$. By Lemma 2.11, for all $i = 1, \ldots, n$, $M'_i = M_i \setminus e_1$ is a multicut of $G$ with $d_{M'_i} = d_{M_i} - 1$. Let $c' = c - \chi e_1$. By $\alpha_v^{y^*} \geq 1$, there exist $\epsilon_i$ for all $i = 1, \ldots, n$, such that $0 \leq \epsilon_i \leq y^*_i$ and $\sum_{i=1}^n \epsilon_i = 1$. The point $y^1$ defined by:

$$y^1 = y^* + \sum_{i=1}^n (-\epsilon_i \xi_{M_i} + \epsilon_i \xi_{M'_i})$$

is a solution to $D_{(G,c')}$. By definition of $b$, we have:

$$b^\top y^1 = b^\top y^* - h - 1. \quad (9)$$

By minimality assumption (ii), $D_{(G,c')}$ admits an integer optimal solution, say $y^2$. This solution satisfies with equality the constraint (8a) associated with $e_2$ as otherwise $y^2 + \xi_{\delta(v)}$ would be a solution to $D_{(G,c)}$ with cost $b^\top y^2 + b_{\delta(v)} \geq b^\top y^1 + 2h + 1$, contradicting the assumption that $y^*$ is optimal by (9) and $h \geq 1$. Hence, there exists a multicut $\bar{M}$ active for $y^2$ containing $e_2$ but not $e_1$ since $c'_1 + 1 \leq c'_2$. By definition, $\bar{M} \cup e_1$ is a multicut of $G$ of order $d_{\bar{M}} + 1$. Define $y^3 \in \mathbb{Z}^{M_G}$ by:

$$y^3 = y^2 - \chi \bar{M} + \chi (\bar{M} \cup e_1)$$

By definition of $c'$ and $y^2$, the point $y^3$ is an integer solution to $D_{(G,c')}$. Therefore, by (9), since $y^2$ is optimal in $D_{(G,c')}$ and by definition of $y^3$, we have:

$$b^\top y^* = b^\top y^1 + h + 1 \leq b^\top y^2 + h + 1 \leq b^\top y^3.$$ 

Thus, $y^3$ is an integer optimal solution to $D_{(G,c)}$, a contradiction.

Suppose now that $\alpha_v^{y^*} = 0$. First, note that $\delta(v)$ is not an active multicut for $y^*$. Otherwise by Claims 5.2, 5.3 and 5.5, there would be a multicut containing $e_1$ and not $e_2$, say $N_1$, and a multicut containing $e_2$ and not $e_1$, say $N_2$, which are both active for $y^*$. This contradicts Claim 5.4. Hence, by definition of $\alpha_v^{y^*}$, no active multicut contains $\delta(v)$.

By Observation 2.2, if a multicut $M$ contains $e_2$ but not $e_1$, then $M \triangle \delta(v)$ is a multicut with the same order and $b_M = b_{M \triangle \delta(v)}$. Hence, we can define the point $y^4 \in \mathbb{Q}^{M_G}$:

$$y^4_M = \begin{cases} 
0 & \text{if } e_1 \notin M, \\
y^*_M + y^*_M \delta(v) & \text{if } e_1 \notin M \text{ and } e_2 \in M, \\
y^*_M & \text{for all } M \in \mathcal{M}_G, 
\end{cases}$$

for all $M \in \mathcal{M}_G$, otherwise.
which is a solution to $\mathcal{D}(G, \hat{c})$, where $\hat{c}$ is defined by:

$$
\hat{c}_e = \begin{cases} 
  c_{e_1} + c_{e_2} & \text{if } e = e_2, \\
  0 & \text{if } e = e_1, \\
  c_e & \text{otherwise},
\end{cases}
$$

By construction, we have:

$$
b^T y^4 = b^T y^*.
$$

Using the argument given in the proof of Claim 5.2, we deduce that $\mathcal{D}(G, \hat{c})$ admits an integer optimal solution, say $y^5$. Let $\mathcal{S}$ be the family of active multicut for $y^5$ containing $e_2$, where each multicut $M$ appears $y^5_M$ times in $\mathcal{S}$. We have $|\mathcal{S}| > c_{e_2}$ as otherwise $y^5$ would be an integer optimal solution to $\mathcal{D}(G, c)$, a contradiction.

We now construct from $y^5$ an integer solution $y^6$ to $\mathcal{D}(G, c)$ with the same cost by replacing $e_2$ by $e_1$ in some active multicut for $y^5$. More formally, since $|\mathcal{S}| > c_{e_1}$, there exists $\mathcal{S}' \subseteq \mathcal{S}$ with $|\mathcal{S}'| = c_{e_1}$. By Observation 2.2, $M \triangle \delta(v)$ is a multicut of $G$ for all $M \in \mathcal{S}'$ and $b_M = b_{M \triangle \delta(v)}$. Let $y^6 \in \mathbb{Z}^{M\mathcal{S}}$ be the point defined by:

$$
y^6 = y^5 + \sum_{M \in \mathcal{S}'} \left( \xi_{M \triangle \delta(v)} - \xi_M \right)
$$

By construction, we have:

$$
b^T y^6 = b^T y^5.
$$

Note that for each $M \in \mathcal{S}'$, adding $\xi_{M \triangle \delta(v)} - \xi_M$ to a point of $\mathbb{R}^{M\mathcal{S}}$ increases (resp. decreases) by 1 the left-hand side of the inequality (8a) associated with $e_1$ (resp. $e_2$) while not changing the left-hand side of the inequalities (8a) associated with the edges of $E \setminus \{e_1, e_2\}$. Therefore, by definition of $\hat{c}$, $y^6$ is a solution to $\mathcal{D}(G, c)$. By (11), the optimality of $y^5$, and we have:

$$
b^T y^6 = b^T y^5 \geq b^T y^4 = b^T y^*.
$$

Therefore $y^6$ is an integer optimal solution to $\mathcal{D}(G, c)$, a contradiction. $\square$

Claim 5.6 implies that for each optimal solution $y$ and for each $v \in \hat{V}$ there exists at least one multicut strictly containing $\delta(v)$ that is active for $y$. For the following claims we need to define a subset of optimal solutions to $\mathcal{D}(G, c)$: let $\mathcal{D}_v$ be the set of optimal solutions to $\mathcal{D}(G, c)$ for which $\delta(v)$ is not active. Note that if $\mathcal{D}_v$ is not empty, then there exists a solution $y$ in $\mathcal{D}_v$ maximizing $\alpha^v_z$ over all $z \in \mathcal{D}_v$.

The following claim presents the structure of a specific optimal solution to $\mathcal{D}(G, c)$ for which $\delta(v)$ is not active.

**Claim 5.7** Let $v \in \hat{V}$ with $\delta(v) = \{e_1, e_2\}$ and let $y^* \in \mathcal{D}_v$ maximize $\alpha^v_z$ over all $z \in \mathcal{D}_v$. Then, there are exactly 3 multicut active for $y^*$ intersecting $\delta(v)$: two bonds $F \cup e_1$ and $F \cup e_2$ and a multicut $F \cup \{e_1, e_2\}$ of order 3, for some $F \subseteq E$.

**Proof** By Claim 5.6, there exists at least one multicut strictly containing $\delta(v)$ which is active for $y^*$, say $M_0$. By definition of $\mathcal{D}_v$, $\delta(v)$ is not active for $y^*$. Hence, by
Claim 5.5, there exists at least one multicut active for \( y^* \) which contains \( e_i \) and not \( \delta(v) \setminus e_i \), for \( i = 1, 2 \). Let \( M_i \) be such a multicut with maximum order.

First, we prove that \( d_{M_0} = 3 \). By definition, \( M_0 = \delta(v, V_2, V_3, \ldots, V_{d_{M_0}}) \). Moreover, by Lemma 2.11 and complementary slackness, the two vertices adjacent to \( v \) belong to two different classes, say \( V_2 \) and \( V_3 \). By contradiction, suppose that \( d_{M_0} \geq 4 \). Then, \( M'_0 = \delta(v \cup V_2 \cup V_3, \ldots, V_{d_{M_0}}) \) is a multicut of order \( d_{M_0} - 2 \). For \( i = 1, 2 \), \( M'_i = M_i \cup \delta(v) \) is a multicut of order \( d_{M_i} + 1 \). Let \( 0 < \varepsilon \leq \min \{ y^*_{M_0}, y^*_{M_1}, y^*_{M_2} \} \).

Then, let \( y' \in \mathbb{R}^{M_G} \) be the point defined by:

\[
y' = y^* - \varepsilon \xi_{M_0} + \varepsilon \xi_{M'_0} + \varepsilon \sum_{i=1,2} (-\xi_{M_i} + \xi_{M'_i}).\]

By construction, \( y' \) is a solution to \( D(G,c) \) with \( b^T y^* = b^T y' \). Hence \( y' \) is an optimal solution, but we have \( \alpha_v y' = \alpha_v y^* + \varepsilon \) because \( \delta(v) \not\subseteq M'_i \) for \( i = 1, 2 \). This contradicts the maximality of \( \alpha_v y^* \). Therefore \( d_{M_0} = 3 \).

Now, we show that \( M_1 \) is a bond. The result for \( M_2 \) holds by symmetry. By contradiction, suppose that \( M_1 = \delta(V_1, \ldots, V_{d_{M_1}}) \) with \( d_{M_1} \geq 3 \). Without loss of generality, we suppose that \( e \in \delta(V_1) \cap \delta(V_2) \). Then, \( M'_1 = \delta(V_1 \cup V_2, \ldots, V_{d_{M_1}}) \) is a multicut of order \( d_{M_1} - 1 \). Moreover, \( M'_2 = M_2 \cup \delta(v) \) is a multicut of order \( d_{M_2} + 1 \). Let \( 0 < \varepsilon \leq \min \{ y^*_{M_1}, y^*_{M_2} \} \) and \( y' \in \mathbb{R}^{M_G} \) be the point defined by:

\[
y' = y^* - \varepsilon \xi_{M_1} + \varepsilon \xi_{M'_1} - \varepsilon \xi_{M_2} + \varepsilon \xi_{M'_2}.
\]

By construction, \( y' \) is a solution to \( D(G,c) \) with \( b^T y^* = b^T y' \). Hence \( y' \) is an optimal solution, but we have \( \alpha_v y' = \alpha_v y^* + \varepsilon \) because \( \delta(v) \not\subseteq M'_2 \). This contradicts the maximality of \( \alpha_v y^* \). Therefore, \( d_{M_1} = d_{M_2} = 2 \).

We now prove that there exists a set \( F \) such that \( M_0 = F \cup \delta(v) \), and \( M_i = F \cup e_i \) for \( i = 1, 2 \). Note that \( M_1 \cup M_2 \) is a multicut so \( y'' = y^* + \varepsilon (\xi_{M_1 \cup M_2} - \xi_{M_1} - \xi_{M_2}) \) is a solution to \( D(G,c) \). The optimality of \( y^* \) implies \( d_{M_1 \cup M_2} \leq 3 \). Since \( M_1 \) and \( M_2 \) are distinct bonds, there exists \( F \subseteq E \setminus \delta(v) \) such that \( M_i = F \cup e_i \), for \( i = 1, 2 \). Finally, let \( N_0 = M_0 \setminus e_2 \) and \( N_1 = M_1 \cup e_2 \). Note that \( \hat{y} = y^* + \varepsilon (\xi_{N_0} - \xi_{M_0} + \xi_{N_1} - \xi_{M_1}) \) is an optimal solution to \( D(G,c) \) for which \( \delta(v) \) is not active. Moreover, \( N_0 \) and \( M_2 \) are bonds active for \( \hat{y} \) since \( d_{M_0} = 3 \). This implies that \( N_0 = F \cup e_1 \), and hence \( M_0 = F \cup \delta(v) \).

This implies that \( M_0, M_1, \) and \( M_2 \) are the only multicuts active for \( y^* \) intersecting \( \delta(v) \). Indeed, if \( M \) is a multicut active for \( y^* \) strictly containing \( \delta(v) \), then repeating the proof above with \( M, M_1, \) and \( M_2 \) shows that there exists \( F' \) such that \( M = F' \cup \delta(v) \), and \( M_i = F' \cup e_i \) for \( i = 1, 2 \). Since \( M_i = F \cup e_i \) for \( i = 1, 2 \), we have \( F' = F \) and hence \( M = M_0 \). A similar argument holds for any multicut active for \( y^* \) intersecting \( \delta(v) \). \( \square \)

Claim 5.8 Let \( v \in \hat{V} \) and \( y^* \) be an optimal solution to \( D(G,c) \). Then,

(i) If \( y^*_{\delta(v)} = 0 \), then \( c_e = 1 \) for all \( e \in \delta(v) \),
(ii) If $y^*_y > 0$, then $\alpha_v^y + y^*_y = 1$, and there exists $e \in \delta(v)$ such that $c_e = 1$.

**Proof**

(i) First suppose that $y^*_y = 0$, then $D_v \neq \emptyset$. Let $y' \in D_v$ maximize $\alpha_v^z$ over all $z \in D_v$. Then, by Claim 5.7, there exist exactly two active multicut for $y'$ containing $e_i$ for $i = 1, 2$. Combining Claims 5.3 and 5.5, and the integrality of $c$, we obtain that $c_{e_i} = 1$ for $i = 1, 2$.

(ii) Let now $y^*_y > 0$. By Claim 5.4, there exists an edge $e \in \delta(v)$ such that all multicuts containing $e$ that are active for $y$ contain $\delta(v)$. Hence, the constraint (8a) associated with $e$ is:

$$c_e \geq \sum_{M : e \in M} y^*_M - y^*_y \sum_{M \in A^*_v} y^*_M = y^*_y + \alpha_v^y. \quad (12)$$

By Claim 5.5, the constraint (8a) associated with $e$ is tight. Thus, $y^*_y + \alpha_v^y = c_e$.

By Claims 5.3 and 5.6 and since $c_e$ is integer, we have $c_e = 1$. \qed

The last three claims of the proof give some structural property of the graph $G$. In particular we focus our attention on the vertices of $\hat{V}$.

**Claim 5.9** Vertices of degree 2 are pairwise nonadjacent.

**Proof** Assume for a contradiction that there exist two adjacent vertices $v_1$ and $v_2$ in $\hat{V}$, and denote $\delta(v_i) = \{e_0, e_i\}$ for $i = 1, 2$.

We prove that $\delta(v_1)$ is active for all optimal solutions to $D_{(G,c)}$, the result for $\delta(v_2)$ is obtained by symmetry. By contradiction, suppose that $D_{v_1} \neq \emptyset$. Among all the solutions $y \in D_{v_1}$, let $y^1$ be one having $\alpha^1_{v_1}$ maximum. Then, by Claim 5.7, the three multicut active for $y^1$ intersecting $\delta(v_1)$ are $M_0 = F \cup \delta(v_1)$, $B_0 = F \cup e_0$, and $B_1 = F \cup e_1$, where $B_1$ are bonds for $i = 0, 1$, and $F \subseteq E \setminus \delta(v_1)$ contains no nonempty multicut. By Claim 5.6 on $v_2$, there exists a multicut $M$ active for $y^1$ strictly containing $\delta(v_2)$. By $\delta(v_1) \cap \delta(v_2) = e_0$, $M$ intersects $\delta(v_1)$. Since $d_M \geq 3$, Claim 5.7 for $v_1$ implies $M = M_0$, $F = \{e_2\}$, and $B_0 = \delta(v_2)$.

As $y^1_{\delta(v_1)} = 0$, by Statement (i) of Claim 5.8, $c_{e_0} = c_{e_1} = 1$. By Claim 5.5, the constraints associated with $e_0$ and $e_1$ are tight. Since $A^1_{v_1} = \{0\}$ by Claim 5.7, we have:

$$c_{e_i} = y^1_{M_0} + y^1_{B_i} = 1 \quad \text{for } i = 0, 1. \quad (13)$$

Let $\{M_1, \ldots, M_n\}$ be the set of active multicut for $y^1$ such that $M_i \cap \{e_0, e_1, e_2\} = e_2$, for $i = 1, \ldots, n$. By Claim 5.5, the constraint (8a) associated with $e_2$ is tight, hence, using (13):

$$c_{e_2} = y^1_{M_0} + y^1_{B_0} + y^1_{B_1} + \sum_{i=1}^n y^1_{M_i} = 1 + y^1_{B_0} + \sum_{i=1}^n y^1_{M_i}. \quad (14)$$
By Claim 5.3, $B_0$ active for $y^1$, and $c_{e_2} \in \mathbb{Z}$, we have $\{M_1, \ldots, M_n\} \neq \emptyset$ and $c_{e_2} \geq 2$. Thus, combining (13) and (14), we have:

$$\sum_{i=1}^{n} y^1_{M_i} = c_{e_2} - 1 - y^1_{B_0} \geq y^1_{M_0}. \quad (15)$$

Then, there exist $\epsilon_1, \ldots, \epsilon_n$ such that $0 \leq \epsilon_i \leq y^1_{M_i}$ for $i = 1, \ldots, n$, and

$$\sum_{i=1}^{n} \epsilon_i = y^1_{M_0}. \quad (16)$$

For $i = 1, \ldots, n$, $M_i \cup e_0$ is a multicut with order $d_{M_i} + 1$, hence we can consider the following solution to $D_{(G,c)}$:

$$y^2 = y^1 - \left(y^1_{M_0} \xi_{M_0} + \sum_{i=1}^{n} \epsilon_i \xi_{M_i}\right) + \left(y^1_{M_0} \xi_{M_0 \setminus e_0} + \sum_{i=1}^{n} \epsilon_i \xi_{M_i \cup e_0}\right). \quad (16)$$

We have $b^\top y^1 = b^\top y^2$, but $\alpha_{\delta(v_1)} y^2 = 0$, a contradiction to Claim 5.6. Therefore $\delta(v) \neq \emptyset$, and by symmetry we deduce that both $\delta(v_1)$ and $\delta(v_2)$ are active for all optimal solutions to $D_{(G,c)}$.

By Claim 5.4, for every optimal solution $y$ to $D_{(G,c)}$ and every multicut $M$ of $G$, if $M$ is active for $y$ and contains $e_i$ for some $i \in \{1, 2\}$, then $e_0 \in M$.

Let $y^*$ be the optimal solution to $D_{(G,c)}$ maximizing $\alpha_{\delta(v)} y$ over all $y$ solutions to $D_{(G,c)}$. We have $A^y_{v_2} \subseteq A^y_{v_1}$ and all the multicuts in $A^y_{v_2}$ have order at most 3. Otherwise, let $M \in A^y_{v_2} \setminus A^y_{v_1}$ (resp. $M \in A^y_{v_2}$ such that $d_M \geq 4$), and $0 < \epsilon \leq \min\{y^*_M, y^*_\delta(v_1)\}$. The solution

$$y^3 = y^* - \epsilon(\xi_{M} + \xi_{\delta(v_1)}) + \epsilon(\xi_{M \setminus e_2} + \xi_{\delta(v_1) \cup e_2})$$

is optimal, but $\alpha_{\delta(v_1)} y^3 = \alpha_{\delta(v_1)} y^* + \epsilon$ by the choice of $M$, a contradiction to the maximality of $\alpha_{\delta(v_1)} y^*$. Thus, $\tilde{M} = \{e_0, e_1, e_2\}$ is the only multicut in $A^y_{v_2}$.

Let $\{N_1, \ldots, N_m\}$ be the set of active multicuts for $y^*$ such that $N_i \cap \{e_0, e_1, e_2\} = e_0$ for $i = 1, \ldots, m$. The constraint associated with $e_0$ is tight by Claim 5.5, hence, by $A^y_{v_2} \subseteq A^y_{v_1}$, we have:

$$c_{e_0} = \alpha_{\delta(v_1)} y^* + y^*_{\delta(v_1)} + y^*_{\delta(v_2)} + \sum_{i=1}^{m} y^*_{N_i}. \quad (17)$$
By Statement (ii) of Claim 5.8 applied to $v_1$, we have $y_{\delta(v_1)}^* + \alpha_{v_1} y^* = 1$, and so:

$$c_{e_0} = 1 + y_{\delta(v_2)}^* + \sum_{i=1}^m y_{N_i}^*. \tag{18}$$

By $A_{v_2}^* = \{ \bar{M} \}$ and Statement (ii) of Claim 5.8 applied to $v_2$, we have $y_{\delta(v_2)}^* + y_{\bar{M}}^* = 1$, hence:

$$c_{e_0} = 2 - y_{\bar{M}}^* + \sum_{i=1}^m y_{N_i}^*. \tag{19}$$

Since $c_{e_0}$ is integer and since $y_{\bar{M}}^* < 1$ by Claim 5.3, by (19), we have:

$$\sum_{i=1}^m y_{N_i}^* \geq y_{\bar{M}}^*. \tag{20}$$

Hence, let $\lambda_1, \ldots, \lambda_m$ be such that $0 \leq \lambda_i \leq y_{N_i}^*$ for $i = 1, \ldots, m$, and $\sum_{i=1}^m \lambda_i = y_{\bar{M}}^*$. Note that $\delta(v_2) = \bar{M} \setminus e_1$. Then, the point

$$y^5 = y^* - \left( y_{\bar{M}}^* \xi_{\bar{M}} + \sum_{i=1}^m \lambda_i \xi_{N_i} \right) + \left( y_{\bar{M}}^* \xi_{\delta(v_2)} + \sum_{i=1}^m \lambda_i \xi_{N_i \cup e_1} \right)$$

is a solution to $D_{(G,c)}$, and it is optimal by definition of $b$. Moreover,

$$y_{\delta(v_2)}^5 = y_{\bar{M}}^* + y_{\delta(v_2)}^* = 1,$$

a contradiction to Claim 5.3.

The following claim forbids a circuit of length 3 to contain a vertex of $\hat{V}$.

**Claim 5.10** No circuit of length 3 contains a vertex of degree 2.

**Proof** Assume for a contradiction that in $G$ there exist a vertex $v \in \hat{V}$ and a circuit $\{e_1, e_2, e_3\}$ such that $\delta(v) = \{e_1, e_2\}$. By Lemma 2.1, a multicut contains $e_3$ only if it intersects $\delta(v)$. On the other hand, by Lemma 2.11 and complementary slackness, each multicut strictly containing $\delta(v)$ and active for an optimal solution contains $e_3$. Thus, for every optimal solution $y$ to $D_{(G,c)}$, we have:

$$\sum_{M : e_3 \in M} y_M = \sum_{M : e_1 \in M, M \neq \delta(v)} y_M + \sum_{M : e_2 \in M, M \neq \delta(v)} y_M - \alpha_v^y. \tag{21}$$

Let $y^*$ be an optimal solution to $D_{(G,c)}$. By the constraint (8a) associated with $e_3$, (21), and Claim 5.5, we have:

$$c_{e_3} \geq \sum_{M : e_3 \in M} y_M^* = c_{e_1} + c_{e_2} - 2 y_{\delta(v)}^* - \alpha_v^y. \tag{22}$$
By Claim 5.6 and Statement (ii) of Claim 5.8, we have $2y_{\delta(v)}^* + \alpha_v^* < 2$. Thus, by (22) and $c_{e_3} \in \mathbb{Z}$, we have $c_{e_3} \geq c_{e_1} + c_{e_2} - 1$.

Define $G' = G \setminus e_3$ and $c' = c|E(\setminus e_3)|$. Note that for each multicut $M \in \mathcal{M}_G$, $M \setminus e_3$ is a multicut of $G'$ with order at least $d_M$. Hence, $y^*$ induces a solution to $\mathcal{D}(G', c')$ of cost at least $b^T y^*$. By minimality assumption (i), there exists an integer optimal solution $y'$ to $\mathcal{D}(G', c')$, and we have $b^T y' \geq b^T y^*$.

Let $\mathcal{M}_1$ (resp. $\mathcal{M}_2$) be the set of multicuts $M = \delta(V_1, \ldots, V_{d_M})$ of $G'$ active for $y'$ such that the endpoints of $e_3$ belong (resp. do not belong) to a same $V_i$ for some $i \in \{1, \ldots, d_M\}$. For each $M \in \mathcal{M}_1$ (resp. $M \in \mathcal{M}_2$), $M$ (resp. $M \cup e_3$) is a multicut of $G$ with the same order. Hence,

$$y'' = \sum_{M \in \mathcal{M}_1} y'_M \xi_M + \sum_{M \in \mathcal{M}_2} y'_M \xi_{M \cup e_3}$$

is a point of $\mathbb{Z}^{|\mathcal{M}_G|}$ with $b^T y'' = b^T y'$. Thus, $b^T y'' \geq b^T y^*$, and $y''$ is not a solution to $\mathcal{D}(G, c)$. By definition, $y''$ respects every constraint of $\mathcal{D}(G, c)$ but the constraint (8a) associated with $e_3$.

By definition of $y''$, we have:

$$\sum_{M : x \in M, M \neq \delta(v)} y''_M = \sum_{M : x \in M, M \neq \delta(v)} y''_M + \sum_{M : e_2 \in M, M \neq \delta(v)} y''_M - \alpha_v^* y''_\delta(v) \leq c_{e_1} + c_{e_2} - \alpha_v^* - 2y_{\delta(v)}^*.$$

(23)

Therefore, by $y''$ violating the constraint (8a) associated with $e_3$, (23), Statement (ii) of Claim 5.8, and the inequalities (8a) associated with $e_1$ and $e_2$, we have:

$$c_{e_3} < \sum_{M : e_3 \in M} y''_M = \sum_{M : x \in M} y''_M + \sum_{M : e_2 \in M} y''_M - \alpha_v^* y''_\delta(v) - 2y_{\delta(v)}^* \leq c_{e_1} + c_{e_2} - \alpha_v^* - 2y_{\delta(v)}^*.$$

(24)

Thus, by (22), we have $\alpha_v^* + 2y_{\delta(v)}^* < \alpha_v^* + 2y_{\delta(v)}^* < 2$. By $c_{e_3} \geq c_{e_2} + c_{e_1} - 1$, the integrality of $y''$, and (24), we have $\alpha_v^* = y''_{\delta(v)} = 0$, and so $c_{e_3} = c_{e_1} + c_{e_2} - 1$. Hence, by the integrality of $y''$ and equation (23):

$$c_{e_3} + 1 = \sum_{M : e_3 \in M} y''_M = \sum_{M : x \in M} y''_M + \sum_{M : e_2 \in M} y''_M = c_{e_1} + c_{e_2}.$$

(25)

For $i = 1, 2$, since $c_{e_i} \geq 1$, there exists a multicut $M_i$ active for $y''$ such that $M_i \cap \delta(v) = e_i$.

We claim that the constraint (8a) associated with $e_3$ is not tight for $y^*$. By $c_{e_3} = c_{e_1} + c_{e_2} - 1$, (22), and Claim 5.6, $\delta(v)$ is active for $y^*$. Hence, by Statement (ii) of Claim 5.8, we have:

$$\alpha_v^* + y_{\delta(v)}^* = 1.$$

(26)

Hence, by (21) and Claim 5.5, (26), (25), and $\delta(v)$ active for $y^*$, we have:

$$\sum_{M : e_3 \in M} y^*_M = c_{e_1} + c_{e_2} - 1 - y_{\delta(v)}^* = c_{e_3} - y_{\delta(v)}^* < c_{e_3}.$$

(27)
The point \( y'' \) respects all the constraints of \( D_{(G,c)} \) except the one associated with \( e_3 \), and this constraint is not tight for \( y^* \). Therefore, there exists \( 0 < \lambda < 1 \) such that

\[
\tilde{y} = \lambda y^* + (1 - \lambda)y''
\]

is a solution to \( D_{(G,c)} \). Moreover, \( \tilde{y} \) is optimal because \( b^\top y^* \leq b^\top y'' \).

All multicuts active for at least one between \( y^* \) and \( y'' \) are active for \( \tilde{y} \). Since \( \delta(v) \) is active for \( y^* \) and \( M_1, M_2 \) are active for \( y'' \), the three multicuts \( M_1, M_2, \) and \( \delta(v) \) are active for \( \tilde{y} \), a contradiction to Claim 5.4.

**Claim 5.11** Each circuit of length 4 contains at most one vertex of degree 2.

**Proof** Assume for a contradiction that there exists a circuit \( C = \{e_1, \ldots, e_4\} \) in \( G \) covering two vertices of \( \tilde{V} \), say \( v_1, v_2 \). By Claim 5.9, \( v_1 \) and \( v_2 \) are not adjacent, hence we assume that \( \delta(v_1) = \{e_1, e_2\} \) and \( \delta(v_2) = \{e_3, e_4\} \). Let \( v_3 \) and \( v_4 \) be the remaining vertices of \( C \).

We prove that \( \delta(v_1) \) is active for all optimal solutions to \( D_{(G,c)} \). Indeed, if \( D_{v_1} \neq \emptyset \), then let \( y' \in D_{v_1} \) maximize \( \alpha_{v_1}^y \) over all \( z \in D_{v_1} \). By Statement (ii) of Theorem 2.10, for every multicut \( M \) in \( A_{v_2}^y \), we have \( M = \delta\{v_2, v_3, \ldots, v_{dy}\} \), with \( v_3 \) and \( v_4 \) belonging to different \( V_i \)'s, hence \( M \cap \delta(v_1) \neq \emptyset \). Moreover, \( M \setminus \delta(v_1) \) contains \( \delta(v_2) \), a contradiction to Claim 5.7 applied to \( v_1 \). Exchanging the role of \( v_1 \) and \( v_2 \), we deduce that \( \delta(v_2) \) is active for all optimal solutions to \( D_{(G,c)} \).

Without loss of generality, there exists an optimal solution \( y \) such that \( \alpha_{v_1}^y \geq \alpha_{v_2}^y \). Then, we can build from \( y \) an optimal solution \( y^* \) to \( D_{(G,c)} \) such that \( A_{v_2}^{y^*} \subseteq A_{v_1}^y \). Indeed, suppose \( A_{v_2}^y \setminus A_{v_1}^y = \{M_1, \ldots, M_n\} \). Then, since \( \alpha_{v_1}^y \geq \alpha_{v_2}^y \), there exist \( N_1, \ldots, N_m \in A_{v_1} \setminus A_{v_2}^y \) such that:

\[
\sum_{i=1}^{n} y_{M_i} \leq \sum_{j=1}^{m} y_{N_j} .
\]

Hence, there exist \( \epsilon_1, \ldots, \epsilon_m \) such that \( 0 \leq \epsilon_j \leq y_{N_j} \), for \( j = 1, \ldots, m \), and

\[
\sum_{j=1}^{m} \epsilon_j = \sum_{i=1}^{n} y_{M_i} .
\]

By Statement (ii) of Theorem 2.10 and complementary slackness, \( v_3 \) and \( v_4 \) belong to different classes of \( N_j \) for each \( j = 1, \ldots, m \), implying that \( N_j \cap \delta(v_2) \neq \emptyset \). Moreover, since \( N_j \notin A_{v_2}^y \), we have \( |N_j \cap \delta(v_2)| = 1 \), for all \( j = 1, \ldots, m \). Furthermore, since \( \delta(v_2) \) is active for \( y \), by Claim 5.4, there exists an edge in \( \delta(v_2) \), say \( e_3 \), such that \( N_j \cap \delta(v_2) = e_3 \) for all \( j = 1, \ldots, m \). Therefore, the point

\[
y^* = y - \left( \sum_{i=1}^{n} y_{M_i} \xi_{M_i} - \sum_{i=1}^{n} y_{M_i} \xi_{M_i \setminus e_4} \right) + \left( \sum_{j=1}^{m} \epsilon_j \xi_{N_j \cup e_4} - \sum_{j=1}^{m} \epsilon_j \xi_{N_j} \right)
\]

is a solution to \( D_{(G,c)} \).
is a solution to $D_{(G,c)}$ with $b^Ty^* = b^Ty$ and $A_{v_1}^{x*} \subseteq A_{v_1}^{y*}$. Let $A_{v_2}^{y*} = \{M'_1, \ldots, M'_p\}$. For $i = 1, \ldots, p$, since $M'_i \in A_{v_1}^{x*}$, Statement (ii) of Theorem 2.10 implies $M'_i = \delta(v_1, v_2, V'_1, V'_2, \ldots, V'_{d_{M'_i}})$, where $V'_3$ and $V'_4$ contain respectively $v_3$ and $v_4$. Then, $M''_i = \delta(v_1, v_2 \cup V'_3, V'_4, \ldots, V'_{d_{M''_i}})$ is a multicut of order $d_{M'_i} - 2$ for $i = 1, \ldots, p$. Since $\delta(v_2)$ is active for $y^*$, by Statement (ii) of Claim 5.8, we have $\alpha_{v_2}^x + y^*_b(v_2) = 1$. Then, the point $y^1 \in \mathbb{Q}^{\mathcal{M}_G'}$ defined by:

$$y^1 = y^* - \left( y^*_{\delta(v_2)} \xi_{\delta(v_2)} + \sum_{i=1}^p y^*_{M'_i} \xi_{M'_i} \right) + \left( \sum_{i=1}^p y^*_{M'_i} \xi_{M''_i} \right),$$

is a solution to $D_{(G,c')}$, where $c' = c - \chi_{\delta(v_2)}$.

By $d_{M''_i} = d_{M'_i} - 2$ for all $i = 1, \ldots, p$, and $\alpha_{v_2}^x + y^*_b(v_2) = 1$, we have:

$$b^Ty^1 = b^Ty^* - \alpha_{v_2}^x(2h + 2) - y^*_{\delta(v_2)}(2h + 1) = b^Ty^* - (2h + 1) - \alpha_{v_2}^x. \quad (31)$$

By minimality assumption (ii), $D_{(G,c')}$ admits an integer optimal solution, say $y^2$. The point $y^3 \in \mathbb{Z}^{\mathcal{M}_G}$ defined by $y^3 = y^2 + \xi_{\delta(v_2)}$ is a solution to $D_{(G,c)}$ such that:

$$b^Ty^3 = b^Ty^2 + 2h + 1. \quad (32)$$

Therefore, by (31), the optimality of $y^2$, and (32), we have:

$$b^Ty^* = b^Ty^1 + 2h + 1 + \alpha_{v_2}^x \leq b^Ty^2 + 2h + 1 + \alpha_{v_2}^x = b^Ty^3 + \alpha_{v_2}^x. \quad (33)$$

By integrality of $P_{2h+1}(G)$ and duality, we have $b^Ty^* \in \mathbb{Z}$. Furthermore, $y^3$ is integer by construction, so $b^Ty^3 \in \mathbb{Z}$. Then, by (33) and Claim 5.6, we have $b^Ty^* \leq b^Ty^3$, and so $y^3$ is an integer optimal solution to $D_{(G,c)}$, a contradiction. □

Claims 5.1, 5.9, 5.10, 5.11 and Lemma 2.3 imply that $G$ is not series–parallel, a contradiction. □

The box-TDIness of $P_k(G)$ and the TDIness of System (2) give the following result.

**Corollary 5.2** System (2) is box-TDI if and only if $G$ is series–parallel.

**Proof** By Theorem 5.1, when $G$ is not series–parallel, System (2) is not TDI. Whenever $G$ is series–parallel, $P_k(G)$ is box-TDI by Theorem 3.2 and System (2) is TDI by Theorem 5.1. □

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