Enlargement of Calderbank Shor Steane quantum codes

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Abstract

It is shown that a classical error correcting code \( C = [n, k, d] \) which contains its dual, \( C^\perp \subseteq C \), and which can be enlarged to \( C' = [n, k' > k + 1, d'] \), can be converted into a quantum code of parameters \([n, k + k' - n, \min(d, \lceil 3d'/2 \rceil)]\). This is a generalisation of a previous construction, it enables many new codes of good efficiency to be discovered. Examples based on classical Bose Chaudhuri Hocquenghem (BCH) codes are discussed.

keywords Quantum error correction, BCH code, CSS code

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Quantum information theory is rapidly becoming a well-established discipline. It shares many of the concepts of classical information theory but involves new subtleties arising from the nature of quantum mechanics [1]. Among the central concepts in common between classical and quantum information is that of error correction, and the error correcting code. Quantum error correcting codes have progressed from their initial discovery [2, 3] and the first general descriptions [4, 3, 5] to broader analyses of the physical principles [8, 10, 11, 13] and various code constructions [11, 13, 17, 9, 12, 16, 14, 15]. A thorough discussion of the principles of quantum coding theory is offered in [6], and many example codes are given, together with a tabulation of codes and bounds on the minimum distance for codeword length $n$ up to $n = 30$ quantum bits.

For larger $n$ there is less progress, and only a few general code constructions are known. The first important quantum code construction is that of [4, 3, 5], the resulting codes are commonly referred to as Calderbank Shor Steane (CSS) codes. It can be shown that efficient CSS codes exist as $n \to \infty$, but on the other hand these codes are not the most efficient possible. I will present here a method which permits most CSS codes to be enlarged, without an attendant reduction in the minimum distance of the code. The resulting codes are therefore more efficient than CSS codes. The examples I will give are found to be among the most efficient quantum codes known, and enable some of the bounds in [6] to be tightened. The code construction is essentially the same as that described for Reed-Muller codes in [16], the new feature is to understand how the method works and thus prove that it remains successful for a much wider class of code. After this some relevant theory of Bose Chaudhuri Hocquenghem (BCH) codes [18, 19, 21] will be given and used to construct a table of example quantum codes built by the new method. The codes are additive and pure in the nomenclature of [13]. A pure additive code is nondegenerate in the nomenclature of [13].

1 Quantum coding

Following [13], the notation $[[n, k, d]]$ is used to refer to a quantum error correcting code for $n$ qubits having $2^k$ codewords and minimum distance $d$. Such a code enables the quantum information to be restored after any set of up to $\lfloor (d-1)/2 \rfloor$ qubits has undergone errors. In addition, when $d$ is even, $d/2$ errors can be detected. We restrict attention to the ‘worst case’ that any defecting qubit (ie any qubit undergoing an unknown interaction) might change state in a completely unknown
way, so all the error processes \(X, Z\) and \(Y = XZ\) must be correctable \([5, 7, 8, 13]\).

A quantum error correcting code is an eigenspace of a commutative subgroup of the group \(E\) of tensor products of Pauli matrices. The commutativity condition can be expressed \([11, 13, 6, 16]\)

\[
H_x \cdot H^T_z + H_z \cdot H^T_x = 0.
\] (1)

where \(H_x\) and \(H_z\) are \((n - k \times n)\) binary matrices which together form the stabilizer \(\mathcal{H} = (H_x | H_z)\). All vectors \((u_x | u_z)\) in the code (where \(u_x\) and \(u_z\) are \(n\)-bit strings) satisfy \(H_x \cdot u_z + H_z \cdot u_x = 0\). These are generated by the generator \(\mathcal{G} = (G_x | G_z)\) which therefore must satisfy

\[
H_x \cdot G^T_z + H_z \cdot G^T_x = 0.
\] (2)

In other words \(\mathcal{H}\) may be obtained from \(\mathcal{G}\) by swapping the \(X\) and \(Z\) parts, and extracting the dual of the resulting \((n + k) \times 2n\) binary matrix. The rows of \(G_x\) and \(G_z\) have length \(n\), and the number of rows is \(n + k\).

The weight of a vector \((u_x | u_z)\) is the Hamming weight of the bitwise OR of \(u_x\) with \(u_z\). The minimum distance \(d\) of the code \(\mathcal{C}\) is the largest weight such that there are no vectors of weight \(< d\) in \(\mathcal{C} \setminus \mathcal{C}^\perp\), where the dual is with respect to the inner product \(((u_x | u_z), (v_x | v_z)) \equiv u_x \cdot v_z + u_z \cdot v_x\). A pure code has furthermore no vectors of weight \(< d\) in \(\mathcal{C}\), apart from the zero vector.

The CSS code construction \([4, 5]\) is to take classical codes \(C_1\) and \(C_2\) with \(C_1^\perp \subseteq C_2\), and form

\[
\mathcal{G} = \begin{pmatrix} G_1 & 0 \\ 0 & G_2 \end{pmatrix}, \quad \mathcal{H} = \begin{pmatrix} H_2 & 0 \\ 0 & H_1 \end{pmatrix}.
\] (3)

where \(G_i\) and \(H_i\) are the classical generator and check matrices. The dual condition \(C_1^\perp \subseteq C_2\) ensures that \(H_1 \cdot H^T_2 = H_2 \cdot H^T_1 = 0\) and therefore the commutativity condition \([11]\) is satisfied. If \(C_1 = [n, k_1, d_1]\) and \(C_2 = [n, k_2, d_2]\) then the minimum distance of the quantum code is \(\min(d_1, d_2)\) and the number of rows in \(\mathcal{G}\) is \(k_1 + k_2\), leading to quantum code parameters \([n, k_1 + k_2 - n, \min(d_1, d_2)]\).

An interesting subset of CSS codes is that given by the above construction starting from a classical \([n, k, d]\) which contains its dual, leading to a quantum \([n, 2k - n, d]\) code.
2 New code construction

I will present the new construction by stating and proving the following.

**Theorem 1.** Given a classical binary error correcting code $C = [n, k, d]$ which contains its dual, $C^\perp \subseteq C$, and which can be enlarged to $C' = [n, k' > k + 1, d']$, a pure quantum code of parameters $[[n, k + k' - n, \min(d, \lceil 3d'/2 \rceil)]]$ can be constructed.

**Proof.** The generator for the quantum code is

$$G = \begin{pmatrix} D & AD \\ G & 0 \\ 0 & G \end{pmatrix}, \tag{4}$$

where $G$ generates the classical code $C$, and $G$ and $D$ together generate $C'$, as does $G$ and $AD$ together (we will choose $A$ such that $D$ and $AD$ generate the same set).

The stabilizer is

$$H = \begin{pmatrix} \tilde{A}B & B \\ H' & 0 \\ 0 & H' \end{pmatrix}, \tag{5}$$

where $H'$ checks the code $C'$, so has $n - k'$ rows, $\{H', B\}$ checks the code $C$, so $B$ has $k' - k$ rows, and

$$\tilde{A} = BD^T \left(A^T\right)^{-1} \left(BD^T\right)^{-1} \tag{6}$$

From the dual conditions specified in the theorem, $H'H'^T = 0$ and $H'B^T = 0$ so the commutativity condition (1) is satisfied. The definition of $\tilde{A}$ ensures we have the correct stabilizer since

$$\tilde{A}B(AD)^T = BD^T. \tag{7}$$

Since the number of rows in the generator is $k + k'$, the dimension of the quantum code is $k + k' - n$. It remains to prove that the minimum distance is $\min(d, \lceil 3d'/2 \rceil)$.

We choose $A$ such that $D$ and $AD$ generate the same set. Therefore for any vector $(u|v)$ generated by $(D|AD)$, either $u = v$ or $\text{wt}(u + v) \geq d'$. We choose the map $A$ such that $u = v$ never occurs (a fixed point free map). This can be achieved...
as long as \( D \) has more than one row, by, for example, the map

\[
A = \begin{pmatrix}
0100\ldots0 \\
0010\ldots0 \\
0001\ldots0 \\
\vdots \\
0000\ldots1 \\
1100\ldots0
\end{pmatrix}.
\] (8)

To complete the proof we will show that for any non-zero vector \((u|v)\) generated by \( G \), \( \mathrm{wt}(u|v) \geq \min(d, 3d'/2) \) (and therefore \( \mathrm{wt}(u|v) \geq \min(d, \lceil 3d'/2 \rceil) \)).

For the non-zero vector \((u|v)\), if either \( \mathrm{wt}(u) \geq d \) or \( \mathrm{wt}(v) \geq d \) then \( \mathrm{wt}(u|v) \geq d \), so the conditions of the theorem are satisfied. The only remaining vectors are those for which both \( \mathrm{wt}(u) < d \) and \( \mathrm{wt}(v) < d \). Now, \( \mathrm{wt}(u) \) can only be less than \( d \) if \( D \) is involved in the generation of \( u \), and \( \mathrm{wt}(v) \) can only be less than \( d \) if \( AD \) is involved in the generation of \( v \), since \( G \) on its own generates a binary code of minimum distance \( d \). However, since the map \( A \) is fixed-point free, and using the fact that \( D \) and \( AD \) generate the same set, the binary vector \( u + v \) is not zero and is a member of a distance \( d' \) code, therefore \( \mathrm{wt}(u + v) \geq d' \). We thus have the conditions \{ \( \mathrm{wt}(u) \geq d' \), \( \mathrm{wt}(v) \geq d' \), \( \mathrm{wt}(u + v) \geq d' \} \). These are sufficient to imply that \( \mathrm{wt}(u|v) \geq 3d'/2 \). For, if \( u \) and \( v \) overlap in \( p \) places, then \( \mathrm{wt}(u + v) = \mathrm{wt}(u) - p + \mathrm{wt}(v) - p \) and \( \mathrm{wt}(u|v) = \mathrm{wt}(u) + \mathrm{wt}(v) - p = \frac{(\mathrm{wt}(u) + \mathrm{wt}(v) + \mathrm{wt}(u + v))/2 \geq 3d'/2} {2} \). This completes the proof.

The above construction was applied to Reed-Muller codes in [16]. These codes are not very efficient (they have small \( k/n \) for given \( n, d \)) but they have the advantage of being easily decoded. A large group of classical codes which combine good efficiency with ease of decoding are the BCH codes. They include Reed Solomon codes as a subset. I will now derive a set of quantum error correcting codes from binary BCH codes using the above construction, combined with some simple BCH coding theory.

3 Application to binary BCH codes

Properties of BCH codes are discussed and proved in, for example, [21]. A binary BCH code of designed distance \( \delta \) is a cyclic code of length \( n \) over GF(2) with
generator polynomial

\[ g(x) = \text{l.c.m.}\{M^{(b)}(x), M^{(b+1)}(x), \ldots, M^{(b+\delta-2)}(x)\} \]  

where

\[ M^{(s)}(x) = \prod_{i \in C_s} (x - \alpha^i) \],

in which \( \alpha \) is a primitive \( n \)th root of unity over GF(2), and \( C_s \) is a cyclotomic coset mod \( n \) over GF(2), defined by

\[ C_s = \{s, 2s, 4s, \ldots, 2^{m_s}-1s\}, \]

where \( m_s = |C_s| \) is obtained from \( 2^{m_s}s \equiv s \mod n \). The dimension of the code is

\[ k = n - \deg(g(x)). \]

From (9) and (10) this implies

\[ k = n - \sum_{s=\delta+1}^{\delta+b-2} |C_s| \]

where the sum ranges from \( s = b \) to \( s = \delta + b - 2 \) but only includes each distinct cyclotomic coset once. This can also be expressed \( k = n - |I_C| \) where \( I_C = C_b \cup C_{b+1} \cup \cdots \cup C_{b+\delta-2} \) is called the defining set. The minimum distance of the code is \( d \geq \delta \).

The dual of a cyclic code is cyclic. Grassl et al. [20] derive the useful criterion that a cyclic code contains its dual if the union of cyclotomic cosets contributing to \( g(x) \) does not contain both \( C_s \) and \( C_{n-s} \). In other words

\[ \{(n-i) \notin I_C \ \forall i \in I_C \} \Rightarrow C^\perp \subseteq C. \]

### 3.1 Primitive BCH codes

Consider first the BCH codes with \( n = 2^m - 1 \), the so-called primitive BCH codes. In order to find the codes which satisfy the condition (12), we will find the smallest \( s \) such that \( n - r \in C_s \) for some \( r \leq s \). The largest permissible designed distance will then be \( \delta = s \). For even \( m \), the choice \( s = 2^{m/2} - 1 \) gives \( s2^{m/2} = n - s \Rightarrow C_s = C_{n-s} \), so this is an upper bound on \( s \). For odd \( m \), an upper bound is provided by \( s = 2^{(m-1)/2} - 1 \) since then \( s2^{(m-1)/2} = n - (s-1)/2 \). We will show that these upper bounds can be filled, i.e. that no smaller \( s \) leads to \( n - r \in C_s \) for \( r \leq s \).

For \( n = 2^m - 1 \), the elements of the cyclotomic cosets \( C_s \) are largest when \( s \) is one less than a power of 2, \( s = 2^i - 1 \). Specifically, for \( s = 2^i - 1 \) we have \( (s2^i \mod n) > (r2^i \mod n) \forall i < m, r < s \). This is obvious for \( s2^i < n \) and the proof for \( s2^i > n \) is straightforward. The largest element in \( C_s \) (\( s = 2^i - 1 \)) is obtained for the largest \( i \) such that \( s2^i < n \), giving \( \max(C_s) = 2^m - 2^{m-j} = n - r \) where
\[ r = 2^{m-j} - 1 \]. This element \( \max(C_s) = n - r \) is the largest in the defining set \( \mathcal{I}_C \) for a code of designed distance \( \delta = s \), therefore it is only possible for \( \mathcal{I}_C \) to contain both \( i \) and \( n - i \) (for any \( i \)) if it contains \( r \) and \( n - r \), since \( r = 2^{m-j} - 1 \) is the smallest element in its coset, and any other pairs \( i, n - i \) must have \( i > r \). Finally, we have a failure of the condition (12) only if \( r \leq s \), that is \( 2^{m-j} - 1 \leq 2^j - 1 \), therefore \( j \geq \lceil m/2 \rfloor \).

To summarise the above, we have the proved following:

**Lemma:** The primitive binary BCH codes contain their duals if and only if the designed distance satisfies
\[ \delta \leq 2^\lceil m/2 \rceil - 1 \] (13)

Using the code construction of theorem 1, together with this lemma, the list of quantum codes in table 1 is obtained. The further property used is that BCH codes are nested, i.e. codes of smaller distance contain those of larger, which is obvious since the former can be obtained from the latter by deleting parity checks. The first entry for each value of \( n \) uses \( \{ C = \text{extended BCH code} \} \) with \( \{ C' = \text{even weight code} \} \) to obtain a distance 3 quantum code. The codes of larger distance involve only BCH codes, for these a quantum code is obtained both from the unextended and extended versions. The parameters \([n, K, D]\) given in the table are for the extended BCH codes (i.e. extended by an overall parity check). Using unextended codes leads to a further quantum code of parameters \([n - 1, K + 1, D - 1]\), for \( D > 3 \).

### 3.2 Non-primitive BCH codes

When \( n \neq 2^m - 1 \) the cyclotomic cosets mod \( n \) do not have so much structure so in general the only way to find if condition (12) is satisfied is to examine each coset individually.

One way in which the requirement (12) is not met is if \( C_s \) contains both \( i \) and \( -i \) mod \( n \), which implies \( C_s = C_{-s} \), for some \( C_s \subseteq \mathcal{I}_C \). If \( s \) is the smallest element in \( C_s \), then \( i, n - i \in C_s \) if and only \( s, n - s \in C_s \), from which \( s2^j \equiv -s \mod n \) for some \( j < m_s \). Multiplying by \( 2^j \) we have \( s2^{2j} \equiv -s2^j \equiv s \mod n \), therefore \( j = m_s/2 \) and this is only possible for even \( m_s \). Furthermore, since \( m_s \geq 1 \) is a factor of \( m_1, m_s \) can be even only if \( m_1 \) is even. This observation slightly reduces the amount of checking to be done.
The values of \( n \) in the range \( 1 < n \leq 127 \) for which \( C_1 \) does not contain \( n - 1 \) are \( \{ 7,15,21,23,31,35,39,45, 47,49,51,55,63,69,71,73,75,77,79,85,87,89,91,93,95,103, 105,111,115,117,119,121,123,127 \} \). An efficient code is obtained if one or more of the cosets is small, this happens for \( n = 21, 23, 45, 51, 73, 85, 89, 93, 105, 117 \) (not counting primitive codes). Quantum codes obtained from BCH codes with these values of \( n \) are listed in table 2. Further good codes exist in the range \( 127 < n < 511 \) for \( n = 133, 151, 153, 155, 165, 189, 195, 217, 219, 255, 267, 273, 275, 279, 315, 337, 341, 381, 399, 455 \).

4 Efficiency

The code parameters in tables 1 and 2 compare well with the most efficient quantum codes known. For example, the [[22, 5, 6]] code permits some of the lower existence bounds in [3] to be raised, and the [[32, 15, 6]] and [[32, 5, 8]] codes fill lower existence bounds. The [[93, 68, 5]] code is comparable with the [[85, 61, 5]] code quoted in [3], though the [[93, 53, 7]] code is not as good as [[85, 53, 7]] quoted in [3]. Obviously the quantum codes based on BCH codes will be best for primitive BCH codes, so we expect the codes in table 1 rather than table 2 to compare best with other code constructions. Indeed, the distance 3 codes in table 1 are the previously known Hamming codes [[13, 6, 3]] and are optimal.

The quantum codes constructed by theorem 1 have an upper bound on the rate \( K/n = (k + k')/n - 1 \) arising from the upper bound on \( k \) and \( k' \) for binary codes. In the asymptotic limit this bound on the quantum codes is

\[
K/n < R(d/n) + R(2d/3n) - 1,
\]

where \( R(d/n) \) is the maximum rate of a binary \([n, k, d]\) linear code. For example the sphere-packing bound is \( R(d/n) < 1 - H(d/2n) \); the codes we have discussed have parameters lying close to this bound (though in the limit of large \( n \) it is known that BCH codes are no longer efficient). Taking \( R(x) \) equal to the McEliece-Rodemich-Rumsey-Welch upper bound [22], we find \( K/n \geq 0 \) for \( d/n < 0.2197 \) in the limit of large \( n \). This may be compared with \( d/n < 0.1825 \) for CSS codes and the limit \( d/n < 0.308 \) for pure quantum stabilizer codes discussed by Ashikhmin [23].

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| $n$ | $k$ | $k'$ | $d$ | $d'$ | $K$ | $D$ |
|-----|-----|------|-----|------|-----|-----|
| 8   | 4   | 7    | 4   | 2    | 3   | 3   |
| 16  | 11  | 15   | 4   | 2    | 10  | 3   |
| 32  | 26  | 31   | 4   | 2    | 25  | 3   |
| 32  | 21  | 26   | 6   | 4    | 15  | 6   |
| 32  | 16  | 21   | 8   | 6    | 5   | 8   |
| 64  | 57  | 63   | 4   | 2    | 56  | 3   |
| 64  | 51  | 57   | 6   | 4    | 44  | 6   |
| 64  | 45  | 51   | 8   | 6    | 32  | 8   |
| 128 | 120 | 127  | 4   | 2    | 119 | 3   |
| 128 | 113 | 120  | 6   | 4    | 105 | 6   |
| 128 | 106 | 113  | 8   | 6    | 91  | 8   |
| 128 | 99  | 113  | 10  | 6    | 84  | 9   |
| 128 | 92  | 106  | 12  | 8    | 70  | 12  |
| 128 | 85  | 99   | 14  | 10   | 56  | 14  |
| 128 | 78  | 99   | 16  | 10   | 49  | 15  |
| 256 | 247 | 255  | 4   | 2    | 246 | 3   |
| 256 | 239 | 247  | 6   | 4    | 230 | 6   |
| 256 | 231 | 239  | 8   | 6    | 214 | 8   |
| 256 | 223 | 239  | 10  | 6    | 206 | 9   |
| 256 | 215 | 231  | 12  | 8    | 190 | 12  |
| 256 | 207 | 223  | 14  | 10   | 174 | 14  |
| 256 | 199 | 223  | 16  | 10   | 166 | 15  |

Table 1. Parameters $[[n, K, D]]$ of the quantum codes obtained from primitive binary BCH codes, for $n \leq 256$. The BCH codes have been extended by an overall parity check in order to allow the distance 3 quantum code to be obtained by combining a BCH code with the even weight code. For $D > 3$ if the unextended BCH codes are used, a $[[n - 1, K + 1, D - 1]]$ quantum code is obtained.
| $n$ | $k$ | $k'$ | $d$ | $d'$ | $K$ | $D$ |
|-----|-----|-----|-----|-----|-----|-----|
| 22  | 15  | 21  | 4   | 2   | 14  | 3   |
| 22  | 12  | 15  | 6   | 4   | 5   | 6   |
| 46  | 33  | 45  | 4   | 2   | 32  | 3   |
| 46  | 29  | 33  | 6   | 4   | 16  | 6   |
| 52  | 43  | 51  | 4   | 2   | 42  | 3   |
| 74  | 64  | 73  | 4   | 2   | 63  | 3   |
| 74  | 55  | 64  | 6   | 4   | 45  | 4   |
| 74  | 46  | 55  | 10  | 6   | 27  | 9   |
| 86  | 77  | 85  | 4   | 2   | 76  | 3   |
| 86  | 69  | 77  | 6   | 4   | 60  | 6   |
| 90  | 78  | 89  | 4   | 2   | 77  | 3   |
| 90  | 67  | 78  | 6   | 4   | 55  | 6   |
| 90  | 56  | 67  | 10  | 6   | 33  | 9   |
| 90  | 45  | 56  | 12  | 10  | 11  | 12  |
| 94  | 83  | 93  | 4   | 2   | 82  | 3   |
| 94  | 78  | 83  | 6   | 4   | 67  | 6   |
| 94  | 68  | 78  | 8   | 6   | 52  | 8   |
| 94  | 58  | 78  | 10  | 6   | 42  | 9   |
| 94  | 53  | 68  | 12  | 8   | 27  | 12  |
| 106 | 93  | 104 | 4   | 2   | 92  | 3   |
| 106 | 81  | 93  | 6   | 4   | 68  | 6   |
| 106 | 75  | 81  | 8   | 6   | 50  | 8   |
| 106 | 71  | 81  | 10  | 6   | 46  | 9   |
| 118 | 105 | 117 | 4   | 2   | 104 | 3   |
| 118 | 93  | 105 | 6   | 4   | 80  | 6   |
| 118 | 81  | 93  | 8   | 6   | 56  | 8   |
| 118 | 69  | 93  | 10  | 6   | 44  | 9   |

Table 2. As table 1, but for non-primitive BCH codes with $n < 127$. 

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