Data mining for detecting Bitcoin Ponzi schemes

Massimo Bartoletti, Barbara Pes, Sergio Serusi

University of Cagliari
Cagliari, Italy

Email: {bart, pes, serusisergio}@unica.it

Abstract—Soon after its introduction in 2009, Bitcoin has been adopted by cyber-criminals, which rely on its pseudonymity to implement virtually untraceable scams. One of the typical scams that operate on Bitcoin are the so-called Ponzi schemes. These are fraudulent investments which repay users with the funds invested by new users that join the scheme, and implode when it is no longer possible to find new investments. Despite being illegal in many countries, Ponzi schemes are now proliferating on Bitcoin, and they keep alluring new victims, who are plundered of millions of dollars. We apply data mining techniques to detect Bitcoin addresses related to Ponzi schemes. Our starting point is a dataset of features of real-world Ponzi schemes, that we construct by analysing, on the Bitcoin blockchain, the transactions used to perform the scams. We use this dataset to experiment with various machine learning algorithms, and we assess their effectiveness through standard validation protocols and performance metrics. The best of the classifiers we have experimented can identify most of the Ponzi schemes in the dataset, with a low number of false positives.

Index Terms—Bitcoin, data mining, fraud detection

I. INTRODUCTION

Bitcoin [1], [2] is decentralized cryptocurrency, which allows secure transfers of money — the bitcoins — without the intermediation of trusted authorities. All transfers of bitcoins are recorded on the blockchain, an immutable public ledger of transactions maintained by a peer-to-peer network through a distributed consensus protocol.

Users can send and receive bitcoins without revealing their true identity: rather, they use pseudonyms (called addresses), which may even be generated fresh for each transaction. Although several approaches to de-anonymize addresses have been proposed [3], [4], [5], [6], [7], specular attempts to strengthen the anonymity of Bitcoin [8], [9], [6], [10], [11] reinforce the perception that criminal activities on Bitcoin are easy to implement, and hard to detect.

Besides classic criminal activities like ransomware [7], [12], [13] and money laundering [14], [15], Bitcoin is currently being used as a payment infrastructure for Ponzi schemes [16]. These are financial frauds disguised as “high-yield” investment programs: actually, a Ponzi scheme repays users only with the funds invested by new users that join the scheme, and so it implodes when it is no longer possible to find new investments [17], [18].

Despite many victims are perfectly aware of their fraudulent nature, and of the fact that they are illegal in many countries, Bitcoin-based Ponzi schemes are proliferating. A recent study [19] inspects the posts on bitcointalk.org (a popular discussion forum on Bitcoin), finding more than 1800 Ponzi schemes from June 2011 to November 2016. Estimating the economic impact of Bitcoin-based Ponzi schemes is more difficult, due to the lack of datasets of Ponzi-related Bitcoin addresses: a conservative estimate for the period from September 2013 to September 2014 shows that Ponzi schemes operated through Bitcoin have gathered more than 7 millions USD. The absence of suitable prevention and intervention policies leads us to believe that many other thousands of victims have been cheated since then, and plundered of millions of USD.

Most of the existing approaches to the analysis of Bitcoin scams require a laborious initial phase of manual or semi-automated search on the web [14], [15], [16], [19], [20], [21], [22] in order to collect Bitcoin addresses involved in the scam. Only after this phase it is possible to automatize the analysis, e.g. to quantify the impact of the scam by inspecting the associated transactions on the blockchain. However, these approaches are ineffective when the scam addresses are not publicly available, e.g. because they are communicated privately to registered users, or published only through the deep web or the dark web. In these cases it would be desirable to have tools that automatically search the Bitcoin blockchain for suspect behaviours, and identify the addresses associated to fraudulent activities.

Given the ever-increasing volumes of data to be managed (∼300 millions of transactions, and several millions of distinct addresses) data mining techniques have become almost imperative for automatically extracting meaningful patterns for fraud detection. Outside the cryptocurrency realm, several works in the literature have explored these techniques with data from credit card operations, either in a supervised setting (which requires a set of labelled observations from the past) [23], [24] or through unsupervised approaches (which look for anomalous data occurrences or outliers) [25], [26]. However, despite an increasing amount of research in the field, practical implementations are rarely reported, as recently pointed out in [27]. Furthermore, the scarcity of publicly available datasets leaves unanswered many questions about which is the best strategy to deal with specific real-world scenarios [28].

The extension of existing fraud detection methods to cybercrime analysis in Bitcoin is an almost unexplored field. A few attempts have been recently made to detect anomalies in the Bitcoin transaction network by unsupervised learning approaches [29], [30], but no work exists, to the best of our knowledge, that investigates how to learn detection models for specific types of scams (such as Ponzi schemes).
Contributions: We investigate data mining techniques to automatically detect and quantify Bitcoin Ponzi schemes, following the supervised learning approach.

In the absence of publicly available datasets, our first step is to retrieve from the web a collection of Bitcoin addresses related to Ponzi schemes. To this purpose we manually search the main discussion forums on Bitcoin (e.g., Reddit and bitcointalk.org) for advertisements of “high-yield” investment programs, that inevitably hide Ponzi schemes. Then, we visit the websites through which Ponzi schemes are operated (possibly recovering old snapshots through Internet Archive), hunting for their Bitcoin addresses. We expand our collection through a semi-automatic visit of the websites that are linked to Bitcoin addresses on blockchain.info/tags. Following this methodology, we collect 32 Bitcoin addresses which gather deposits from investors of Ponzi schemes.

In many cases, Ponzi schemes use multiple addresses: actually, some of them provide the deposit address only upon registration, generating a fresh address for each new user. In order to retrieve some of these addresses, we apply a clustering procedure on the addresses in our collection, using the “multi-input” heuristic [3]. By analyzing the obtained clusters, we find that 19 out of 32 Ponzi schemes in our collection use more then one address, for a total of 1211 addresses. Overall, these clusters have received deposits for ~10 millions USD.

We then devise a set of features that can be useful to characterise Ponzi schemes. These features range from simple statistics on the transactions to/from the clusters (e.g., overall transferred value, ratio between incoming and outgoing transactions, etc.) to more complex ones, like measures of inequality of the transferred values (e.g., Gini coefficients), and measures of the activity of the scheme (e.g., lifetime, average delay between incoming and outgoing transactions, maximum number of daily transactions, etc.). We extract from the Bitcoin blockchain the transactions of the clusters of addresses in our collection, and we compute a dataset of features, which we make publicly available. We complete this dataset with the features of 6400 randomly-chosen addresses.

We use this dataset to experiment with various supervised learning algorithms, in order to automatically detect Ponzi schemes. We formalise the detection model as a binary classification problem, where the task is to distinguish between ‘Ponzi’ and ‘non-Ponzi’ class instances. One of the most critical challenges we had to face is the class imbalance problem, which is commonly encountered in fraud detection systems [31]. In a supervised learning setting, as the one here considered, this problem occurs when one class is very rare compared to the other(s), thus making hard to discover robust patterns for the minority class (like “finding a needle in a haystack”). Indeed, classifiers are usually designed to minimize the total number of classification errors, and tend to be overwhelmed by the majority class.

In fraud detection applications, as in many domains with imbalanced class distributions, a correct classification of the rare class (i.e., the ‘Ponzi’ class in our problem) is far more important than a correct classification of the majority class (i.e., the ‘non-Ponzi’ class). The underlying assumption is that the cost of misclassifying a fraudulent case is much higher than the cost of misclassifying a legitimate case (as the latter error can be corrected a posteriori through a further analysis). In this work we experiment with the two main approaches proposed in the literature, i.e. sampling-based approaches [32] and cost-sensitive approaches [33].

A number of experiments across different settings resulted in a detection model with good performance, which is finally applied, with promising results, to an independent set of data. The supervised method Random Forest proved to be the most effective and most versatile one. In our dataset, containing the features of 6432 clusters of addresses (proportion of one fraud to 200 not fraud), Random Forest has obtained a Recall of 0.969 for Ponzi schemes, and it has classified correctly 31 Ponzi schemes out of 32.

In summary, our main contributions are:

1) a public dataset of addresses and features of Bitcoin Ponzi schemes (goo.gl/ToCho7);
2) an open-source tool that extracts the dataset from the Bitcoin blockchain (github.com/bitcoinponzi);
3) a systematic evaluation and comparison of different learning strategies for classifying Bitcoin Ponzi schemes;
4) the evaluation of the best classifier (among those we have experimented with) on an independent dataset, which manages to identify most of the Ponzi schemes in the dataset, with a low number of false positives;
5) an estimate of which are the most discriminating features for detecting Ponzi schemes on Bitcoin.

The rest of this paper is organized as follows. Section II gives a minimalistic introduction to Bitcoin. Section III illustrates our methodology for collecting addresses of Ponzi schemes, and for constructing a dataset of Ponzi-related features. Section IV compares the effectiveness of various learning strategies. Finally, Section V draws some conclusions.

II. BITCOIN IN A NUTSHELL

In this section we give a short introduction to Bitcoin, focussing on the notion that are needed later on for our technical development.

Bitcoin is a peer-to-peer infrastructure which allow users to transfer currency — the bitcoins (฿). Each Bitcoin user owns one or more pairs of asymmetric cryptographic keys: public keys uniquely identify the user addresses, while private keys are used to authorize payments. Transfers of bitcoins are described by transactions. The log of all transactions, is recorded on a public, immutable data structure (the blockchain), determining the balance of each address. Users can receive bitcoins through different addresses: typically, addresses are generated fresh for each transaction, to improve privacy.

The Bitcoin network is populated by a large set of nodes, called miners, which collect transactions from users, and are in charge of appending the valid ones to the blockchain. To this purpose, each miner keeps a local copy of the blockchain, and a set of transactions received by users. Appending a new
block of transactions to the blockchain requires miners to solve a moderately-hard cryptographic puzzle, which involves the transactions in the new block. The difficulty of the puzzle is adjusted dynamically to ensure that the average mining rate is of 1 block every 10 minutes. The miner which solves the puzzle before the others receives a reward in newly generated bitcoins (through the so-called coinbase transactions), and a fee for each transaction included in the new block.

To explain how transactions work, we consider the example in Figure 1, which graphically represents three transactions. Each transaction has four (indexed) fields: in, wit, out, and val. The in field (for input) contains a reference to the transaction output to redeem. The wit field contains a piece of information called witness, discussed below. The field out contains an output script: intuitively, this is a predicate on one or more arguments, the actual values of which are provided by the witness of the redeeming transaction. Finally, the field val determines the amount of bitcoins to be transferred.

Consider first the transaction $T$ in our example, where we neglect the fields in and wit (e.g., we could assume that $T$ is a coinbase transaction). The output at index 1 allows to transfer 1฿ to user $A$: namely, the output script $\text{versig}_A(x)$ verifies the signature $x$ of $A$ on the redeeming transaction. Similarly, the output at index 2 allows to transfer 2฿ to $B$. Assume that the blockchain contains $T$, and that both its outputs are not redeemed by subsequent transactions.

Transaction $T_A$ has one input, represented as the pair $(T, 1)$, meaning that it wants to redeem the output at index 1 of transaction $T$. To do so, $T_A$ carries in its witness a signature of $A$, which is computed on the whole transaction $T_A$ (except for the wit field itself). This witness makes the output script $\text{versig}_A(x)$ in $T$ evaluate to true. Therefore, when $A$ appends $T_A$ to the blockchain, it redeems the first output of $T$, making available $0.9$ ฿ for $B$, and keeping $0.1$ ฿ for herself.

The transaction $T_B$ has two inputs, meaning that it wants to simultaneously redeem the second output of $T$ and the first output of $T_A$. Since these outputs are still unspent, and the witnesses in $T_B$ satisfy the corresponding output scripts, then $B$ can append $T_B$ to the blockchain, making available $2.5$ ฿ to user $C$. The difference of $0.4$ ฿ between the the sum of the values of all the inputs of $T_B$ and the sum of the values of its outputs is paid as a fee to the miners.

Bitcoin transactions may be more general than the ones illustrated by the previous example. For instance, the output script is a program in a (not Turing-complete) scripting language, featuring a limited set of logic, arithmetic, and cryptographic operators. Transactions can also specify time constraints on when they can be appended to the blockchain, and also on when the redeeming transactions can be appended. We omit a detailed presentation of these advanced features, since they are not required in the following sections.

### III. Dataset construction

The first step of our work is to collect Bitcoin addresses through which Ponzi schemes receive money from investors (Section III-A). We apply a clustering algorithm to the collected addresses, finding that some schemes use wallets of hundreds of addresses (Section III-B). We then devise a set of features that are relevant to the classification of Ponzi schemes (Section III-C), and we compute the values of these features on our clustered addresses, obtaining a dataset that we use in Section IV to train classifiers (Section III-D).

#### A. Collection of Bitcoin addresses used by Ponzi schemes

We perform a manual search on Reddit and bitcointalk.org, the main discussion forums on Bitcoin. In particular, we focus on the subforum Gambling: Investor-based games of bitcointalk.org, where fraudsters are used to advertise Ponzi schemes as “high-yield investment programs” (HYIP), or as gambling games. Only in a few cases these advertisements explicitly include the Bitcoin address where to deposit money; in all the other cases, to obtain the address we have to visit the websites where the Ponzi schemes are hosted. However, many of these websites are no longer online: in such case we try to recover their snapshots through Internet Archive. For each website (either live or snapshot), we manually search its pages to find the deposit addresses (typically, different “investment plans” use different addresses). Some websites only allow registered users to read the deposit address: in these cases, we create an account, providing fake data.

We extend our search by considering all the addresses listed on blockchain.info/tags, a website which allows users to tag Bitcoin addresses. Most of the tagged addresses also contain a link to the website where they are mentioned. We develop a crawler to automatically parse these websites, and rank them according to the number of Ponzi-related words contained in their pages. To this purpose we use a dictionary, containing words like e.g. “Ponzi”, “profit”, “HYIP”, “multiplier”, “investment”, “MLM”. The crawler parses over 1500 websites (related to ∼3500 tags), finding that ∼900 of them contain some Ponzi-related word. However, many of these sites (∼600) are no longer accessible, even through Internet Archive. For the remaining websites, we manually search for deposit addresses, creating fake accounts when needed.

Overall, we find 32 deposit addresses of Ponzi schemes, that we display in Table 1. Note that, while some Ponzi schemes use a single deposit address throughout their lives, some others use multiple addresses, possibly generating a fresh address for each user (or set of users). Address clustering, discussed below, allows us to recover some of these addresses.

| $T$ | $T_A$ | $T_B$ |
|-----|-------|-------|
| in: T | in: (T, 1) | in: (T, 2) |
| wit: ··· | wit: $\text{sig}_A(T_A)$ | wit: $\text{sig}_B(T_B)$ |
| out1(x): $\text{versig}_A(x)$ | out1(x): $\text{versig}_A(x)$ | out1(x): $\text{versig}_B(x)$ |
| val1: 1 Satoshi | val1: 0.9 Satoshi | val1: 0.4 Satoshi |
| out2(y): $\text{versig}_B(y)$ | out2(y): $\text{versig}_A(y)$ | out2(y): $\text{versig}_A(y)$ |
| val2: 2 Satoshi | val2: 0.1 Satoshi | val2: 2.5 Satoshi |

**Figure 1: Three Bitcoin transactions.**
Table I: Collection of addresses of Bitcoin Ponzi schemes.

| Ponzi scheme                | Deposit address                                                                 |
|----------------------------|--------------------------------------------------------------------------------|
| Nanonindustry.com          | 1E5zhZ2I7wvkyMvXUy5xYVSX3Ie4vNYEAGA                                         |
| GrandAgofinance            | 1MzQ7H8Q9dQ86Q5Z2B8GylCzkWv2P3sVjG6                                           |
| Cryptory                   | 1FyedwPKKd1342vwvfPhyV1C1L3bVGLxHzXPK                                         |
| Leancy                     | 15sMMioDzfAnwfdB5qNw6xM8h5nCNx8EwikRvY                                       |
| Minimalist10               | 1U9vywA4Cw7n34yMw5B5bQ126nV2561naf6WgBwB                                        |
| MiniPonziCoin              | 1fZKpiI9M2p9f79w01pxZ6R0nK5zXm3fT1f1d                                        |
| 12cycle                    | 1E5MCTXsOv7nZvpyzv3z1bHD6y9y797q0Q6zP5                                           |
| 10PERCENTBTC               | 1Btc8So98e9m8FICeU5cCnq30eEF8Y7azpnc                                           |
| bcogains                   | 1Puyon8Rl8wz6yZymgea3qyfPphg8Dv5                                               |
| PonziFIO                   | 1ponzi1ujCvBi617Zm7TTH4rAURr1vEe6q                                                |
| LaxoTrade                  | 1LaxoTroy1Ly1B289V9mo5Ag666fUj7bL9                                               |
| OpenPonzi                  | 1BM2w6S5ozWae1L9L9MPLnK8s18InqUSMau                                              |
| BTC-doubler.us             | 1AQp51H22wH2DzrLF48656nOu3b685T183Q22                                         |
| BTC-doubler.com            | 178uAKKxwTsx4T3vKZit6gZLzjyEG2e3                                                   |
| investorbitco.com          | 1CpV3Ae84BvzijfH8hgeZlZLV1111o6Qg                                                |
| Ponzi120                   | 12PvNZ6tawbkCU4YF8f8WhNFC1FtAaK3m5q                                              |
| RockwellPartners           | 139e6GkMGR69F9F69q3q4Y0x8bknBnASlV                                               |
| Twelverized                | 114Ap9C0n37Bv8E8c64884mPw3QeCuoVPv5SL                                               |
| CRYPTOSX2                  | 19ZYlMBJm8xAx8574aQvuD6cT6rK                                                    |
| Ihourbtc, pw               | 1BjsjAHT2Qh5s82x5H6fNz1UItuWhtxxoKbNEN                                            |
| bitcoindoubler.fund        | 1FNgsgBhynmMUFMXMsFeZ2bgaMsS95c                                                   |
| doublebtc.life             | 1zmnw5eWBWpryPv5nS2K9r7ThuXa09c                                                 |
| bitcoincopy.site123.me     | 18A6yd9Tszxzv9y5nTmbabyyrw23qSSKhA                                                |
| bitcoinprofit12            | 1AX1qWvYi12bDZlq9g3ZvpsjBnpx                                                  |
| invest4profit              | 1PZ85eO77eUc5v12gbc7j2x8yc7wtxG                                                 |
| 1getpaid.me                | 1GetP9dxmWx3NK7JYc9wvqR2c82zE                                             |
| Ponzi.io (change)          | 14j9kmgNHlIc7f8K7JYwmjG6756                                            |
| igain.com                  | 1AQcSpP94jP7q3pPX16h78vE668nK3YQF                                               |
| world-btc.online           | 1A886eD6q9XHRBMCMocyK5B9QhpAuz256                                             |
| bestdoubler.eu             | 13NztAxzK5n8cUC2U5PqF7KdX7TDJ/F3Mg7Lh                                          |
| bitcoindoubler.prv.pl      | 18Mkvy7fl4N4z589fjC66Nh8S7m3ZkEinV                                              |

**B. Address clustering**

Many techniques to break the anonymity of Bitcoin users have been proposed in the literature. This is achieved either by grouping together (“clustering”) the addresses controlled by each user [3], [4], [5], or by using observations on the underlying peer-to-peer network [34], [35], or by combining both techniques [36].

Several heuristics for address clustering have been proposed over the years. Besides analysing the shape of the transaction graph, some heuristics also take into consideration the behavior of standard clients [5]. To construct our dataset of Ponzi-related addresses, we use the multi-input heuristic [3], [4], the simplest and most efficient one. The key assumption of this heuristic is that, in a multi-input transaction (like e.g. $T_B$ in Figure 1), all the addresses referred to within the inputs are controlled by the same user. These transactions occur, for instance, when a user A wishes to transfer a certain amount of $v$ to another user B, but none of the transactions in A’s wallet has an unspent output of at least $v$. In this case, to avoid paying multiple transaction fees, A can perform the transfer in a single shot, by putting on the blockchain a multi-input transaction redeemable by B, where the sum of the values redeemed by the inputs is at least $v$. Typical Bitcoin clients implement this by choosing the input transactions from A’s wallet, satisfying the assumption of the multi-input heuristic.

We show in Table II some statistics on the clusters that we obtain after applying the multi-input heuristic to the 32 addresses in our collection. The columns display the size of the clusters, the overall number of transactions (either including or outgoing), and the overall inflow, both in $\text{addr}$ and in USD. To convert the amount of each transaction to USD, we use the average exchange rate on the day of the transaction, obtained from www.coindesk.com/price. Overall, the Ponzi schemes in our collection gathered almost 10 millions USD; the scheme that raised the most is Cryptory, with ~ 4.6 millions USD.

Table II: Top-10 Ponzi schemes by cluster size.

| Ponzi scheme                | #Addr. | #Tx | In ($M) | In ($) |
|----------------------------|-------|-----|---------|-------|
| LaxoTrade                  | 491    | 4798 | 1,580    | 570,106 |
| Cryptory                   | 232    | 222  | 9,439    | 6,658,008 |
| Ihourbtc                   | 180    | 1262 | 36      | 42,668  |
| 120cycle                   | 78     | 224  | 14      | 8265    |
| bitcoindoubler.fund        | 63     | 1143 | 90      | 288,849 |
| world-btc.onlin            | 40     | 302  | 1      | 2060    |
| Ponzi.io                   | 33     | 6311 | 370     | 258,368 |
| bcogains                   | 14     | 789  | 72      | 33,246  |
| 10PERCENTBTC               | 13     | 10,077 | 107      | 42,894  |
| investorbitcoin.com        | 11     | 672  | 312     | 158,569 |

Total (32 schemes) 1211 107,637 17,910 9,509,050

C. Features extraction

We now introduce a set of features, which are relevant for the classification of Bitcoin addresses.

- **The lifetime** of the address, expressed in number of days. This is computed as the difference between the date of the first transaction to the address, and the date of the last transaction from/to the address.
- **The activity days**, i.e. the number of days in which there has been at least a transaction from/to the address.
- **The maximum number of daily transactions from/to the address.**
- **The Gini coefficient of the values transferred to (resp. from) the address.** Gini coefficients are a standard representation of the degree of inequality of wealth: 0 indicates perfect equality, while 100 is perfect inequality [16].
- **The sum of all the values (resp. to/from) the address.**
- **The number of incoming (resp. outgoing) transactions which transfer money to (resp. from) the address.**
- **The ratio between incoming and outgoing transactions from/to the address.**

All the features above are defined “pointwise” on single Bitcoin address. We extend them to clusters in the straightforward way: any feature on a cluster is the composition of the pointwise features on the addresses included in the cluster.
As an additional “componentwise” feature, we consider the number of addresses included in the cluster.

D. Dataset construction

We construct a binary dataset that contains two classes of instances: Ponzi schemes (denoted as $P$) and others (denoted as $nP$). Each instance in the dataset corresponds to a cluster of Bitcoin addresses (computed as shown in Section III-B), and it is represented as a tuple of features (the ones defined in Section III-C, plus the class label $P$ or $nP$). To compute these dataset instances we exploit an open-source tool for custom blockchain analytics [37].

We populate the dataset with 32 instances of the class $P$ (corresponding to our clusters of Ponzi schemes), and with 6400 randomly-chosen instances of the class $nP$ (which are clustered with the multi-input heuristic as well). This strong imbalance between the two classes (approximately, 1 Ponzi instance every 200 instances of non-Ponzi) is needed to properly model the fact that Ponzi-related addresses are extremely rare, compared to non-Ponzi ones. Although 1/200 is still much higher than the expected ratio between Ponzi addresses and non-Ponzi ones, it is a necessary compromise to meaningfully represent also the rare class in the dataset.

IV. DATA MINING FOR PONZI SCHEMES

We formalise the induction of a detection model for Bitcoin Ponzi schemes as a binary classification problem, where the task is to distinguish between ‘Ponzi’ and ‘non-Ponzi’ class instances. The strategies for dealing with the imbalanced distribution of the classes are discussed below (Section IV-A), along with the learning algorithms applied to induce the model (Section IV-B), the performance metrics and the validation protocol (Section IV-C). A number of experiments across different settings (Section IV-D) resulted in a detection model with good performance, which is then applied, with promising results, to an independent set of data (Section IV-E). Finally, we investigate which are the most relevant features, among those in our set, to detect Ponzi schemes (Section IV-F).

A. Class imbalance problem

The class imbalance problem is one of the most critical issues faced by fraud detection systems [31]. In a supervised learning setting, as the one here considered, this problem occurs when one class is very rare compared to the other(s), thus making hard to discover robust patterns for the minority class. Indeed, classifiers are usually designed to minimize the total number of classification errors, and tend to be overwhelmed by the majority class.

In fraud detection applications, as in many domains with imbalanced class distributions, a correct classification of the rare class (i.e., the $P$ class in our problem) is far more important than a correct classification of the majority class (i.e., the $nP$ class), since the cost of misclassifying a fraudulent case is usually higher than the cost of misclassifying a legitimate case (as the latter can be corrected through ex-post analyses).

A number of approaches have been proposed in the literature for handling this problem, including sampling-based approaches [32] and cost-sensitive approaches [33].

Sampling-based approaches: The basic idea is to modify the distribution of instances so that the minority class is adequately represented in the dataset used for model development. The most common sampling technique is random undersampling (RUS), which consists in removing observations at random from the majority class. An alternative approach is random oversampling (ROS), where some of the minority instances are replicated, but with an increased risk of overfitting, particularly with noisy data [28]. Though more sophisticated (and expensive) approaches exist, they have not proved to be superior in severe imbalance settings [38]. Furthermore, the effectiveness of sampling techniques may be dependent on the learning algorithm used (and on the adopted performance measure); as well, the extent of sampling for best performance may be domain-dependent.

Cost-sensitive approaches: Cost-sensitive learning involves the use of a cost matrix which encodes the penalty of classifying instances from one class as another. In a class imbalance setting, where the focus is usually on rare instances, a misclassification of the minority class is penalized more than a misclassification of the majority class. For a given classification model, penalty terms are then used to derive an overall cost which reflects the “weight” of the different types of classification errors, besides their total number. Cost-sensitive classification techniques take this cost matrix into consideration in the training phase, in order to generate the classification model with the lowest overall cost.

In this work we consider both sampling-based and cost-sensitive approaches. Furthermore, we experiment with learners whose inner design can cope, at least to some extent, with imbalanced class distributions, as in the case of the RIPPER algorithm proposed in [39].

B. Classifiers

In the induction stage of our detection model, we exploited RIPPER, Bayes Network and Random Forest classifiers, which are representatives of quite different learning strategies.

RIPPER is a propositional rule learner that relies on a sequential covering logic [39] to extract classification rules directly from training data. Rules are grown in a greedy fashion, starting from empty rule antecedents and repeatedly adding conjuncts in order to maximize the information gain measure. An incremental reduced error pruning technique is used the refine the resulting rules. Since the algorithm is designed to give higher priority to the least frequent class, this approach is particularly suited for dealing with imbalanced classification tasks, as in the case of fraud detection.

Bayes Network is a probabilistic model that represents, in the form of a directed acyclic graph, the relation of conditional dependence among a set of variables (the features
and the target class, in the context of classification problems). Probabilistic parameters are encoded in a set of tables, one for each node of the network, in the form of local conditional distributions of a variable given its parents. Both the network structure and the probability values can be estimated from a training set of labelled instances. Bayesian models have been applied in the context of fraud detection systems, e.g. in [40].

**Random Forest** is an ensemble method that exploits multiple decision trees built from random variants of the same data [41]. Although a single tree may be unstable and overly sensitive to the specific composition of the training set, the aggregation of the predictions made by the individual trees in the forest has been shown to achieve much better performance. Compared to other ensemble approaches, Random Forest is computationally efficient and has proved to be a “best of class” learner in several domains [42], including fraud detection [23].

For the above classifiers, we leverage the implementation provided by the Weka machine learning library [43].

C. Performance measures and validation

To evaluate the performance of our detection models, we rely on best practices from the literature. Specifically, in the context of binary problems with imbalanced class distributions, as in the case here considered, the rare class is denoted as the positive class, while the majority class is denoted as the negative class. The following terminology is then used to describe how the model performs on a given set of test instances: a true positive (resp. negative) is a positive (resp. negative) instance correctly classified by the model; a false negative is a positive instance wrongly classified as negative; a false positive is a negative instance wrongly classified as positive.

Depending on the specific characteristics of the data at hand, different metrics can be used for quantifying the extent to which the model is able to recognize positive and negative instances [43]. Hereafter, TP (resp. TN) refers to the number of true positives (resp. negatives), while FP (resp. FN) refers to the number of false positives (resp. negatives):

- **Accuracy** \( (TP + TN)/(TP + TN + FP + FN) \) is the fraction of test instances whose class is predicted correctly;
- **Specificity** \( (TN/(TN+FP)) \) is the fraction of negative instances classified correctly;
- **Sensitivity** \( (TP/(TP+FN)) \), also called Recall, is the fraction of positive instances classified correctly;
- **Precision** \( (TP/(TP+FP)) \) is the fraction of instances that actually are positive in the group the model has predicted as positive;
- **F-measure** \( (2 \cdot \text{Precision} \cdot \text{Recall}) / (\text{Precision} + \text{Recall}) \) is the harmonic mean between precision and recall;
- **G-mean** \( (\text{Recall} \cdot \text{Specificity})^{0.5} \) is the geometric mean between specificity and recall;
- **AUC** is the area under the Receiver Operating Characteristics (ROC) curve, which shows the trade-off between true positive and false positive rates (the better the model, the closer the area is to 1).

Accuracy is the most common performance metric but, alone, is not suited for evaluating models induced from imbalanced datasets. In a fraud detection context, if 0.5% of the instances are fraudulent (as in our dataset), then a model that predicts every instance as non-fraudulent has an accuracy of 99.5%, even though it fails to detect any of the frauds. In this situation, class-specific metrics (such as specificity, recall and precision) can help to better describe and understand the model behaviour. In particular, recall and precision are widely used in applications where the successful detection of the rare class is considered more interesting or important (as for the ‘Ponzi’ class in our problem).

To avoid building models that maximize one metric at the expense of another, trade-off values such those expressed by F-measure and G-mean are taken into account. In turn, AUC is usually considered more significant than accuracy when comparing the overall performance of different classifiers.

Instead of using a single test set to compute the above metrics, we adopt an iterative cross-validation protocol, which involves splitting the original dataset into \( K \) subsets (folds). At each iteration, one of the folds is retained as test data for evaluating the model performance, while the remaining \( K - 1 \) folds are used as training data for building the model. This procedure is repeated \( K \) times, using each time a different fold as test set, and the results of the \( K \) runs are finally aggregated to obtain TP, TN, FP and FN counts. In our experiments, based on common practise in the literature, we set \( K = 10 \).

D. Results

We now present the results obtained with RIPPER, Bayes Net and Random Forest classifiers presented in Section IV-B. Hereafter, we use the following acronyms: RIP for RIPPER, BN for Bayes Net and RF for Random Forest.

First, we evaluate their performance without applying the sampling-based and the cost-sensitive approaches (see Section IV-A). The results are shown in Figure 2 in the form of confusion matrices, where the row index refers to the actual class, while the column index refers to the predicted class. As we can see, the classification performance is not satisfactory, due to the high number of Ponzi schemes not recognized. Bayes Net classifies correctly the largest number of Ponzi instances (23 out of 32), but the number of false positives (i.e., non-Ponzi classified as Ponzi) is the highest as well.

These results confirm that learning from highly imbalanced datasets is a very difficult task.

As a next step, we explore the effectiveness of the random undersampling approach, which has proved to be useful to deal with datasets where the fraud rate is comparable with the one here considered [23]. Specifically, at each iteration of the cross-validation procedure, we manipulate the training set to reduce the extent of class imbalance: the original proportion of 1 Ponzi instance every 200 instances of non-Ponzi (1:200) is reduced to 1 Ponzi every 40 non-Ponzi (1:40), 1 Ponzi every 20 non-Ponzi (1:20), 1 Ponzi every 10 non-Ponzi (1:10) and 1 Ponzi every 5 non-Ponzi (1:5). Note that we do not modify
the class distribution of the test instances, to not introduce any bias in the final performance estimate.

The results in Figure 3 show that the undersampling approach results in an improved true positive rate. In particular, within the 1:5 setting, Random Forest recognizes the same number of Ponzi as Bayes Net (25 out of 32), but with a significantly lower number of false positives. Bayes Net, indeed, produces too many false positives (266), even more than RIPPER (226). Thus, while improving the true positive rate (and hence the recall metric), the undersampling approach is not quite satisfactory in terms of false positives (that affect the precision metric). This difficulty of achieving an optimal trade-off between recall and precision is a recognized issue in the fraud detection literature [31].

As a further step, we investigate the effectiveness of the cost-sensitive approach. When learning our detection models, we use the cost matrices shown in Figure 4. In the matrix CM5, the cost of committing a false negative error is 5 times larger than the cost of committing a false positive error; it is 10 times larger in CM10, 20 times larger in CM20 and 40 times larger in CM40.

The results achieved by the cost-sensitive classifiers are shown in Figure 5. As we can see from the confusion matrices, RIPPER obtains the same results as in the original setting (Figure 2). This is not surprising, since the algorithm is designed to cope with the rare (i.e., positive) class and turns out to be insensitive to further penalising a wrong classification of the positive instances. In turn, Bayes Net does not seem to take significant advantage of the cost-sensitive approach, which results in 24 true positives (one more than in the original setting), but with an increased number of false positives.

The best results are obtained with the Random Forest classifier. Using the CM5 matrix, it recognizes 25 Ponzi schemes, as Bayes Net in the 1:5 undersampling setting, but with a strong reduction of the false positives (only 13). When penalising more the false negatives, the number of the true positives increases (29 using CM10 and 31 using CM20) and the number of false positives increases in turn, but to an acceptable extent (26 e 77 respectively). Increasing further the cost (CM40) is not beneficial since the number of true positives remains the same but the false positives increase to 132.

In Figure 6, we further detail the performance of the cost-sensitive Random Forest classifier, that has shown to be superior to the other approaches here explored. Different performance metrics are computed as explained in Section IV-C.

In terms of accuracy, which simply expresses the fraction of correctly classified instances (irrespective of their class), the best result is achieved with the CM5 cost matrix. It also ensures the highest true negative rate (specificity) and a good trade-off between recall and precision (in terms of F-measure). However, given the specific characteristics of the considered domain, where the correct classification of Ponzi schemes is of paramount importance, we consider especially relevant the recall value, which is optimised using the CM20 cost matrix. The G-mean value, that expresses a trade-off between recall and specificity, is also better with CM20 and, in this setting, the AUC value is the highest as well.

Taking these considerations into account, the Random Forest model obtained using CM20 can be considered the most effective for detecting Ponzi schemes.

E. Application of the induced model

In this section perform an ex-post validation of the best classifier obtained so far, i.e. Random Forest with CM20. To this purpose we collect other Ponzi schemes by searching the web, with the same methodology of Section III-A. We report their addresses in Table III. Overall, the 20 Ponzi schemes in this collection have gathered more than 15 millions USD, in large part with a single scheme, CryptoSplit (see Table IV).

We then construct an alternative dataset, comprising the features of the clusters of the Ponzi schemes in the new collection, and those of 4000 randomly-chosen Bitcoin addresses not in the original dataset. By applying the Random Forest classifier with CM20 to the alternative dataset, we obtain the following confusion matrix:

|   | P    | nP   |
|---|------|------|
| P | 18   | 2    |
| nP| 81   | 3919 |

Notably, the classifier recognizes 18 Ponzi schemes out of 20, producing 81 false positives. The 2 Ponzi schemes not recognized by the classifier are marked with $\times$ in Table IV.

F. Ranking and evaluation of features

We now study which features, among those listed in Section III-C, are more relevant for the classification of Ponzi schemes. To this purpose we exploit the feature selection functionality of Weka [43], which implements several methods for ranking features. Among them, we apply some univariate methods (Information Gain, Gain Ratio, Symmetrical Uncertainty, and OneR), and the multivariate method ReliefF.

Among the 32 features included in our dataset, we consider those with the highest number of occurrences in the first positions of the rankings, and thus can be considered as the most discriminating ones. These features are the following:
Figure 2: Confusion matrices for RIPPER, Bayes Net and Random Forest.

Figure 3: Confusion matrices for RIPPER, Bayes Net and Random Forest across different undersampling in training data.

Figure 4: Cost matrices: CM5, CM10, CM20, and CM40.

Figure 5: Confusion matrices of RIPPER, Bayes Net and Random Forest across different cost-matrices.

| Random Forest | Accuracy | Recall | Specificity | F-measure | Precision | G-mean | AUC |
|---------------|----------|--------|-------------|-----------|-----------|--------|-----|
| CM5: Using Cost 5 | .997 | .781 | .998 | .714 | .658 | .883 | .890 |
| CM10: Using Cost 10 | .995 | .906 | .995 | .667 | .527 | .949 | .951 |
| CM20: Using Cost 20 | .988 | .969 | .987 | .443 | .287 | .978 | .978 |
| CM40: Using Cost 40 | .979 | .969 | .979 | .318 | .190 | .973 | .974 |

Figure 6: Performance of Random Forest across different cost-matrices.
(i) the Gini coefficient of the outgoing values; (ii) the ratio between incoming and total transactions; (iii) the average and standard deviation of the outgoing values; (iv) the number of different addresses who have transferred money to the cluster, and subsequently received money from it; (v) the lifetime of the cluster, and the number of activity days.

V. CONCLUSIONS

Criminal activities that accept payments in bitcoins damage the reputation of Bitcoin, and eventually may be detrimental to the diffusion of cryptocurrencies for legitimate uses. However, since all currency transfers are recorded on a public ledger, surveillance authorities can analyse them, trying to detect anomalous or suspect behaviours.

Despite the transparency of the blockchain, tracking illicit financial flows is a challenging problem, for several reasons. First, many illicit activities involve hundreds, or even thousands, of transactions — thus making manual inspection impracticable. For instance, the Ponzi schemes in our dataset use ~1400 transactions on average: this number is sufficiently large to discourage any attempt at manual inspection. Second, the number of illicit flows is overwhelmed by that of legitimate ones, making the task of surveillance authorities similar to finding “the needle in the haystack”. Another difficulty is that smart cyber-criminals exploit techniques to make analysing their activities more difficult, e.g. by using mixing services to hide the actual provenance of illegal money. All these observations highlight the pressing need for automated techniques to detect illegal activities on cryptocurrencies.

In this work we have proposed an automatic analysis of Ponzi schemes on Bitcoin, based on supervised learning algorithms. Ponzi schemes are a classic fraud masqueraded as “high-yield” investment schemes. However, in a Ponzi scheme the investors are repaid only with the funds invested by new users, hence eventually the scheme implodes, as at a certain point it will no longer be possible to find new investments.

After a preliminary phase of manual search, we have identified 32 Bitcoin addresses used by Ponzi schemes. Address clustering allowed us to extend our collection to 1211 addresses, which overall received investments for ~ 10 millions USD. We have devised a set of features of clusters, that we have used to create a dataset containing the features of all the addresses of Ponzi schemes, and those of other 6400 randomly-chosen addresses. We experimented with data mining tools to evaluate different supervised learning strategies. The best classifier we have found correctly classifies 31 Ponzi schemes out of 32, producing ~ 1% of false positives.

An obvious extension of this work is to apply this classifier to all the addresses in the Bitcoin blockchain. This kind of analysis poses serious efficiency issues, since several dozens of millions of distinct Bitcoin addresses have been used so far. Although the number of false positives is quite low (comparable to that of other successful approaches in the fraud detection literature [23]), automated techniques to check the false positives are in order. To this purpose one could exploit auxiliary information sources, e.g. web discussion forums, and the IP addresses collected by monitoring the traffic on the Bitcoin network.

Our classifier can also be used to detect Ponzi schemes implemented over other cryptocurrencies, like e.g. Ethereum. To this purpose we could exploit public datasets of Ethereum-based Ponzi schemes [44], which collect addresses and other relevant data of 152 Ponzi schemes. In the case of Ethereum, the precision of the classifier could be improved by exploiting more specific features, like e.g. the distribution of gain among users, and the correlation between the timings of inflows and outflows observed in [44].

The approach we have followed in this work can be exploited for the detection of other cryptocurrency-based frauds besides Ponzi schemes, like e.g. ransomware, money laundering, etc. This would require, as in our case, a preliminary phase of dataset construction. The dataset could take into account, besides the features used in Section III-C, further features that better capture specific behaviours of the fraud under analysis.

A relevant question is what interventions can be devised after an illegal activity has been detected. The ex-post sanitization of fraudulent activities is hampered by the current fungibility of the Bitcoin currency. This means that Bitcoin users and exchanges are not selective in which bitcoins to accept, and which ones to reject. Hence, even if we set up risk scores for Bitcoin transactions as proposed in [21], e.g. by marking as “bad” all the bitcoins flowing out a Ponzi scheme, it would not be possible to take countermeasures to the use of “bad” bitcoins until they leave the Bitcoin ecosystem through an exchange service.

REFERENCES

[1] S. Nakamoto, “Bitcoin: a peer-to-peer electronic cash system,” https://bitcoin.org/bitcoin.pdf, 2008.
