The Kalmanson Complex

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Abstract

Let \( X \) be a finite set of cardinality \( n \). The Kalmanson complex \( \mathcal{K}_n \) is the simplicial complex whose vertices are non-trivial \( X \)-splits, and whose facets are maximal circular split systems over \( X \). In this paper we examine \( \mathcal{K}_n \) from three perspectives. In addition to the \( T \)-theoretic description, we show that \( \mathcal{K}_n \) has a geometric realization as the Kalmanson conditions on a finite metric. A third description arises in terms of binary matrices which possess the circular ones property. We prove the equivalence of these three definitions. This leads to a simplified proof of the well-known equivalence between Kalmanson and circular decomposable metrics, as well as a partial description of the \( f \)-vector of \( \mathcal{K}_n \).

Keywords: circular split system; Kalmanson matrix; \( f \)-vector; simplicial complex; phylogenetic network; forbidden substructure

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1 Introduction

A phylogenetic tree is a connected, acyclic graph which presents the common evolutionary history of a group of species (taxa). A phylogenetic network generalizes this structure by allowing for the presence of cycles. Phylogenetic networks have become a popular means of conveying recombination, horizontal transfer and other reticulate events which cannot be represented by a tree [3, 11].

A particularly simple and well-known form of phylogenetic network is the split network. As explained in [6], split networks are mathematically founded on $T$-theory (cf. [10]) and express the so-called circular decomposition of a finite metric [4, 8]. The necessary conditions for such a decomposition are a set of linear inequalities, so that the space of circular decomposable metrics possesses polyhedral structure.

In this paper we investigate that structure. Permuting these inequalities produces a set of polyhedra whose union contains all circular decomposable metrics. Since the polyhedra intersect along faces, the resulting face lattice forms a simplicial complex. (We call this the Kalmanson complex after [12], who derived the original inequalities while studying certain tractable instances of the traveling salesman problem.)

Abstracting away from the underlying geometry, we show how combinatorially isomorphic objects can be derived in terms of either circular split systems, or binary matrices which possess the consecutive ones property [5]. Our main result is to prove the equivalence of these structures by exhibiting order-preserving bijections between their face lattices (Theorems 2 and 12). This in turn leads to a new proof of the equivalence of Kalmanson and circular decomposable metrics (Corollary 9). This is a known result [7, 8], but our proof has the benefit of being extremely simple, relying only on basic concepts from polyhedral geometry.

We then use these findings to study the $f$-vector of Kalmanson complex. We relate the problem of enumerating its faces to a counting problem on certain classes of binary matrices, and exploit a structure theorem of [18] to obtain a new result on the number of triangles contained (Theorem 19). Even in the simplest non-trivial case, this counting problem is seen to possess considerable complexity, and we leave a more general method of counting the faces of the Kalmanson complex as an interesting open problem.

The paper is organized as follows. Section 2 begins with some preliminaries from $T$-theory which enable us to define the complex abstractly in terms of split systems. In Section 3 we review the Kalmanson conditions. These are a set of inequality restrictions on a finite metric which, when satisfied, allow the traveling salesman problem to be solved in constant time. We show that these inequalities give a geometric realization of the Kalmanson complex. In Section 4 we study the consecutive ones property for binary matrices. This property is shown to be equivalent to the circularity property for split systems discussed above, giving us a third description of the complex...
in terms of equivalence classes of binary matrices. In Section 5 we use these three ways of viewing the Kalmanson complex to enumerate some of its faces, thus giving a partial characterization of its $f$-vector. Finally, in Section 6 we offer some concluding remarks.

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2 Definitions

We begin with some basic concepts from $T$-theory. For a full introduction, see [4, 10]. Throughout the paper, $X$ is a finite set of cardinality $n \geq 4$. An $X$-split is a bipartition of $X$; that is, $S = \{A, B\}$ is an $X$-split if $A \cap B = \emptyset$ and $A \cup B = X$. (When the meaning is obvious, we will simply call $S$ a split.) $A$ and $B$ are called the blocks of $S$, and the size of $S$ is defined as $\text{size}(S) := \min\{|A|, |B|\}$. $S$ is non-trivial if $\text{size}(S) > 1$ and minimal if $\text{size}(S) = 2$.

Let $S(X)$ be the set of non-trivial $X$-splits. A split system $S \subset S(X)$ is a set of splits. [4] introduced the concept of a circular split system$^1$

**Definition 1.** A split system $S$ is circular if there is a permutation $\sigma \in S_n$ such that for each split $S = \{A, B\} \in S$ there exists $i, j \in [n]$ such that

$$S = \left\{ \left\{ x_{\sigma(i)}, x_{\sigma(i+1)}, \ldots, x_{\sigma(j)}, x_{\sigma(j+1)}, \ldots, x_{\sigma(i)}, x_{\sigma(i+1)} \right\}, \left\{ x_{\sigma(j)}, x_{\sigma(j+1)}, \ldots, x_{\sigma(i)}, x_{\sigma(i+1)} \right\} \right\}$$

where $\bar{i}$ denotes $i \pmod{n}$.

Circular split systems have a simple geometric interpretation: they are obtained by labeling the edges of a regular $n$-gon, and connecting its edges with diagonals to form splits (Figure 1). From this we see that a circular split system contains at most $\binom{n}{2}$ distinct splits. A partial converse also holds: a weakly compatible (cf. Section 5) split system containing $\binom{n}{2}$ splits is circular [4].

From Definition 1 we see the set of circular split systems is closed under the operations of taking subsets (any subset of a circular split system is circular) and forming intersections. Hence, it is a simplicial complex. This complex is our main object of study.

$^1$Circular split systems are sometimes referred to in the literature as cyclic split systems.
Definition 2. The Kalmanson complex is the simplicial complex whose vertices are $X$-splits, and whose facets are maximal circular split systems.

Clearly this complex is unique up to the cardinality of $X$. Henceforth we write $\mathcal{K}_n$ to denote the Kalmanson complex over a base set of cardinality $n$.

3 Geometry of $\mathcal{K}_n$

Throughout this section, we let $D = (d_{ij})_{i,j \in [n]}$ be a symmetric, non-negative matrix with zeros along the diagonal. We refer to matrices possessing this property as distance matrices.

Definition 3. Let $D = (d_{ij})$ be a distance matrix and let $S_n$ denote the symmetric group on $n$ letters. The traveling salesman problem (TSP) over $D$ is

$$\min_{\sigma \in S_n} \left( \sum_{i=1}^{n-1} d_{\sigma(i)\sigma(i+1)} + d_{\sigma(n)\sigma(1)} \right)$$

For general $D$, it is well-known that the TSP is NP-hard. However, some special cases have lower complexity. In particular, [12] showed that if $D$ satisfies a certain set of linear inequalities, then the TSP over $D$ possesses a trivial solution.

Theorem 1 ([12]). Let $D$ be a distance matrix. If

$$\max(d_{ij} + d_{kl}, d_{il} + d_{jk}) \leq d_{ik} + d_{jl} \text{ for all } 1 \leq i < j < k < l \leq n$$  

(1)
then the identity permutation solves the TSP over \( D \).

The inequalities \([1]\) are referred to as the Kalmanson conditions, and a matrix which satisfies them is a Kalmanson matrix (or simply Kalmanson.)

It may be that \( D \) does not satisfy \([1]\), but that some permutation of the rows and columns of \( D \) does. In this case we say that \( D \) is a permuted Kalmanson matrix. Since permuting \( D \) amounts to simply relabeling the underlying distance or cost data, this operation preserves the structure of the problem. \([8]\) give an \( O(n^2) \) recognition algorithm for permuted Kalmanson matrices, so we say the TSP is polynomial time-solvable for this class.

Geometrically, \([1]\) comprises a finite intersection of closed half-spaces: a polyhedron. Given a polyhedron \( P \subset \mathbb{R}^k \) and a hyperplane \( H \subset \mathbb{R}^k \), we say \( H \) supports \( P \) if \( H \cap P \neq \emptyset \) and \( P \) is completely contained in one of the closed half-spaces defined by \( H \). \( F \subset P \) is a face of \( P \) if \( F = P \cap H \) for some supporting hyperplane \( H \) of \( P \). The face lattice of \( P \) is the poset of faces of \( P \) ordered by set inclusion.

Recall that a set of polyhedra which intersect along faces is called a polyhedral fan. Permuting the indices in \([1]\) generates a polyhedral fan which we denote \( \mathcal{P}_n \).

**Example 1.** For \( n = 4 \), \( \mathcal{P}_n \) is the union of three polyhedra obtained by permuting the indices in \([1]\): \( \mathcal{P}_4 := (d_{ij})_{i,j \in [4]} \) such that

\[
\begin{align*}
    d_{12} + d_{34} &\leq d_{14} + d_{23} \\
    d_{13} + d_{24} &\leq d_{14} + d_{23}
\end{align*}
\]

or

\[
\begin{align*}
    d_{13} + d_{24} &\leq d_{12} + d_{34} \\
    d_{14} + d_{23} &\leq d_{12} + d_{34}
\end{align*}
\]

or

\[
\begin{align*}
    d_{14} + d_{23} &\leq d_{13} + d_{24} \\
    d_{12} + d_{34} &\leq d_{13} + d_{24}
\end{align*}
\]

Collectively, these define the region of \( \mathbb{R}^{(4)} \) containing all \( 4 \times 4 \) permuted Kalmanson matrices.

### 3.1 Equivalence of \( \mathcal{P}_n \) and \( \mathcal{K}_n \)

The main claim of this section is that \( \mathcal{P}_n \) is a geometric realization of \( \mathcal{K}_n \) in the sense that they are combinatorially equivalent.

**Theorem 2.** The face lattices of \( \mathcal{K}_n \) and \( \mathcal{P}_n \) are isomorphic as posets.

The remainder of this subsection is devoted to proving the theorem by finding an inclusion-preserving bijection between the faces of these two sets.
In [9] it is shown that the polyhedron defined by (1) decomposes into an $n$-dimensional lineality space and a pointed cone of dimension $\binom{n}{2} - n$. We are interested in the structure of the latter since it encapsulates the combinatorial data embodied by the polyhedron. The authors give an explicit description of the extreme rays of this cone.

**Example 2.** For $n = 5$, the rays of the standard Kalmanson polyhedron are

\[
V^{(2)} = \begin{bmatrix}
0 & 0 & 1 & 1 & 1 \\
0 & 0 & 1 & 1 & 1 \\
1 & 1 & 0 & 0 & 0 \\
1 & 1 & 0 & 0 & 0 \\
1 & 1 & 0 & 0 & 0
\end{bmatrix}, \quad V^{(3)} = \begin{bmatrix}
0 & 0 & 0 & 1 & 1 \\
0 & 0 & 0 & 1 & 1 \\
1 & 1 & 1 & 0 & 0 \\
1 & 1 & 1 & 0 & 0
\end{bmatrix}
\]

\[
V^{(1,3)} = \begin{bmatrix}
0 & 1 & 1 & 1 & 0 \\
1 & 0 & 0 & 1 & 1 \\
1 & 0 & 0 & 1 & 1 \\
0 & 1 & 1 & 0 & 0 \\
0 & 1 & 1 & 0 & 0
\end{bmatrix}, \quad V^{(1,4)} = \begin{bmatrix}
0 & 1 & 1 & 1 & 0 \\
1 & 0 & 0 & 0 & 1 \\
1 & 0 & 0 & 0 & 1 \\
1 & 0 & 0 & 0 & 1 \\
0 & 1 & 1 & 1 & 0
\end{bmatrix}, \quad V^{(2,4)} = \begin{bmatrix}
0 & 0 & 1 & 1 & 0 \\
0 & 0 & 1 & 1 & 0 \\
1 & 1 & 0 & 0 & 1 \\
1 & 1 & 0 & 0 & 1 \\
0 & 0 & 1 & 1 & 0
\end{bmatrix}
\]

We see that the $V^{(i)}$ and $V^{(i,j)}$ have a structure which is the result of arranging square blocks of zeros along the diagonal. It turns out that these matrices, along with their permutations, encode the non-trivial $X$-splits.

**Definition 4.** Let $S = \{A, B\}$ be an $X$-split. A split metric $\delta_S : X \times X \to \mathbb{R}$ is a function such that

\[
\delta_S(x, y) = \begin{cases} 
0, & \{x, y\} \subset A \text{ or } \{x, y\} \subset B \\
1, & \text{otherwise}
\end{cases}
\]

Split metrics are unique.

**Lemma 3.** Let $S_1, S_2 \in S(X)$. If $\delta_{S_1} = \delta_{S_2}$, then $S_1 = S_2$.

**Proof.** For a split $S$ define $\gamma_S(i) := \{x \in X : \delta_S(1, x) = i\}$. We have $S = \{A, B\} = \{\gamma_S(0), \gamma_S(1)\}$. Hence $S_1 = \{\gamma_{S_1}(0), \gamma_{S_1}(1)\} = \{\gamma_{S_2}(0), \gamma_{S_2}(1)\} = S_2$. \hfill \qed

Returning to the example, let $\Delta_S$ be the (symmetric, $n \times n$) matrix associated to $\delta_S$. The matrices
in Example 2 are obtained from split metrics:

\[
V^{(2)} = \Delta_{12|345} \quad V^{(3)} = \Delta_{123|45} \\
V^{(1,3)} = \Delta_{145|23} \quad V^{(1,4)} = \Delta_{14235} \quad V^{(2,4)} = \Delta_{125|34}
\]

We now formalize this idea.

**Theorem 4** \(^{(9)}\). The space of Kalmanson matrices consists of an \(n\)-dimensional lineality space and an \(n(n - 3)/2\)-dimensional pointed cone. The lineality space is spanned by the matrices \(E^{(i)} = \left(e^{(i)}_{pq}\right), 1 \leq i \leq n\) where

\[
e^{(i)}_{pq} = \begin{cases} 
1, & p = i \text{ xor } q = i \\
0, & \text{otherwise}
\end{cases}
\]

(2)

The pointed cone is ruled by the symmetric matrices \(V^{(i)} = \left(v^{(i)}_{pq}\right), 2 \leq i \leq n - 2\) and \(V^{(i,j)} = \left(v^{(i,j)}_{pq}\right), 1 \leq i \leq n - 3, i + 2 \leq j \leq n - 1\), where

\[
v^{(i)}_{pq} := \begin{cases} 
1, & 1 \leq p \leq i < q \leq n \\
0, & \text{otherwise}
\end{cases}
\]

(3)

\[
v^{(i,j)}_{pq} := \begin{cases} 
1, & 1 \leq p \leq i < q \leq j \text{ or } i < p \leq j < q \leq n \\
0, & \text{otherwise}
\end{cases}
\]

(4)

Now let the symmetric group \(S_n\) act on the set of \(n \times n\) matrices by symmetric permutation of rows and columns: \(\sigma \cdot M = (m_{\sigma(i),\sigma(j)})\) for all \(\sigma \in S_n\) and \(M = (m_{ij})\). The following lemma is immediate from \((3)\) and \((4)\).

**Lemma 5.** \(V^{(i)} = \Delta_{1 \ldots i|i+1 \ldots n}\) and \(V^{(i,j)} = \Delta_{i+1 \ldots j|j+1 \ldots i}\). Additionally, symmetrically permuting \(V^{(i)}\) and \(V^{(i,j)}\) is equivalent to applying the same permutation to the underlying split: \(\sigma \cdot V^{(i)} = \Delta_{\sigma(1) \ldots \sigma(i)|\sigma(i+1) \ldots \sigma(n)}\) (and similarly for \(V^{(i,j)}\)).

Define

\[
\mathcal{V} := \{V^{(i)} : 2 \leq i \leq n - 2\} \cup \{V^{(i,j)} : 1 \leq i \leq n - 3, i + 2 \leq j \leq n - 1\}
\]

\[
\mathcal{R} := \{\sigma \cdot V : \sigma \in S_n, V \in \mathcal{V}\}
\]
Thus $\mathcal{R}$ is the set of vertices (rays) of $\mathcal{P}_n$. Finally, let $T(V)$ be the map which takes a matrix in $\mathcal{R}$ to its corresponding split,

$$T : \mathcal{R} \to \mathcal{S}(X)$$

$$\sigma \cdot V(i) \mapsto \left\{ \{\sigma(1), \ldots, \sigma(i)\}, \{\sigma(i+1), \ldots, \sigma(n)\} \right\}$$

(Note that $T$ is well-defined since for each $i, j \in [n]$ there exists a $\sigma \in S_n$ such that $V(i,j) = \sigma \cdot V(j-i).$)

**Lemma 6.** $T : \mathcal{R} \to \mathcal{S}(X)$ is a bijection.

*Proof.* Injectivity follows from the uniqueness of split metrics. For surjectivity, let $T = \{A, B\} \in \mathcal{S}(X)$ be a split. Let $\sigma \in S_n$ be a permutation such that $\sigma^{-1} \cdot A = \{1, \ldots, |A|\}$. Then $T(\sigma \cdot V(|A|)) = S$. □

It remains to show that $T$ is order-preserving: $T(U) \subset T(V) \iff U \subset V$. This follows from the fact that $T$ maps faces to faces.

**Lemma 7.** $\{T(M_1), T(M_2), \ldots, T(M_k)\}$ is a face of $\mathcal{K}_n$ if and only if $\{M_1, M_2, \ldots, M_k\}$ is a face of $\mathcal{P}_n$.

*Proof.* If $\{T(M_1), T(M_2), \ldots, T(M_k)\}$ is a face of $\mathcal{K}_n$ then it is circular respect to the ordering $(\sigma(1), \ldots, \sigma(n))$ for some $\sigma \in S_n$. Hence $\{M_1, M_2, \ldots, M_k\} \subseteq \sigma \cdot \mathcal{V}$ is a face of $\mathcal{P}_n$.

Conversely, since

$$T(V(i)) = \left\{ \{1, 2, \ldots, i\}, \{i + 1, \ldots, n\} \right\}$$

$$T(V(i,j)) = \left\{ \{1, 2, \ldots, i, j + 1, \ldots, n\}, \{i + 1, \ldots, j\} \right\}$$

the set $S := T(\mathcal{V})$ is a maximal circular split system with the ordering $(1, 2, \ldots, n)$. Therefore the claim is true when the $M_i \in \mathcal{V}$.

Now if $M_1, \ldots, M_k$ is an arbitrary face of $\mathcal{P}_n$, then there is a $\sigma \in S_n$ such that for each $1 \leq i \leq k$ there exists $V_i \in \mathcal{V}$ with $M_i = \sigma \cdot V_i$. Then

$$\{T(M_1), \ldots, T(M_k)\} = \{T(\sigma \cdot V_1), \ldots, T(\sigma \cdot V_k)\}$$

$$= \{\sigma \cdot T(V_1), \ldots, \sigma \cdot T(V_k)\}$$

$$\subseteq \sigma \cdot S$$

is a face of $\mathcal{K}_n$. □
This concludes the proof of Theorem 2.

### 3.2 Circular Decomposability

A metric $\delta$ is circular decomposable if it can be written as the positively-weighted sum of circular split metrics, i.e.

$$\delta = \sum_{S \in \mathcal{C}} \alpha_S \delta_S$$

for some circular split system $\mathcal{C}$ and weights $\alpha_S > 0$. It has been shown in [7] and [8] that $\delta$ is circular decomposable if and only if it satisfies the Kalmanson conditions. Both proofs are non-trivial; [7] relies on a Crofton-type formula for computing distances in metric spaces, while [8] uses a number of results from the theory of metrics over a finite set [4].

The polyhedral characterization of the Kalmanson cone given in [9], coupled with the observations of the preceding section, enable us to establish this equivalence in a new and straightforward way. Indeed, $\delta$ satisfies the Kalmanson conditions iff it is in a permuted Kalmanson cone. By Theorem 4 this occurs iff $\delta$ is a linear combination of the permuted matrices $E^{(i)}, V^{(i)}, V^{(i,j)}$ for some $\sigma \in S_n$:

$$\delta = \sigma \cdot \left( \sum_{i=1}^{n} \alpha_i E^{(i)} + \sum_{i=2}^{n-2} \beta_i V^{(i)} + \sum_{i=1}^{n-3} \sum_{j=i+2}^{n-1} \gamma_{ij} V^{(i,j)} \right) \text{ for } \alpha_i \in \mathbb{R} \text{ and } \beta_i, \gamma_{ij} > 0 \quad (5)$$

By Lemma 8, $V^{(i)}$ and $V^{(i,j)}$ are split metrics, and it is easily seen that the $E^{(i)}$ are split metrics corresponding to trivial splits.

Now, equation (5) is not necessarily a circular decomposition since the $\alpha_i$ can be negative. However, assuming $\delta$ obeys the triangle inequality (which, recall, is not implied by the Kalmanson conditions), the $\alpha_i$ are seen to be non-negative.

**Lemma 8.** Let $\delta$ be written as in (5). Then

$$\delta(i, i+1) + \delta(i+1, i+2) - \delta(i, i+2) = 2\alpha_{i+1}$$

where $i, i + 1, i + 2$ are modulo $n$.

**Proof.** Put $\delta = \delta_\alpha + \delta_\beta + \delta_\gamma$, where

$$\delta_\alpha = \sum_{i=1}^{n} \alpha_i E^{(i)} \quad \delta_\beta = \sum_{i=2}^{n-2} \beta_i V^{(i)} \quad \delta_\gamma = \sum_{i=1}^{n-3} \sum_{j=i+2}^{n-1} \gamma_{ij} V^{(i,j)}$$

For $i, i + 1, i + 2$ modulo $n$. $\square$
From equations (2)−(4) we have
\[
\delta_\alpha(i, j) = \alpha_i + \alpha_j \quad (6)
\]
\[
\delta_\beta(i, j) = \sum_{2 \leq s \leq n-2 \atop i \leq s < j} \beta_s \quad (7)
\]
\[
\delta_\gamma(i, j) = \sum_{i \leq s < j \leq t \atop 1 \leq s \leq n-3} \gamma_{st} + \sum_{s < i \leq t < j \atop 1 \leq s \leq n-3 \atop s+2 \leq t \leq n-1} \gamma_{st} \quad (8)
\]

Hence \(\delta_\alpha(i, i+1) + \delta_\alpha(i+1, i+2) - \delta_\alpha(i, i+2) = 2\alpha_{i+1}\).

For \(\delta_\beta\) we have
\[
\delta_\beta(i, i+1) = \begin{cases} 
0, & i = 1, n - 1 \\
\beta_i, & \text{otherwise}
\end{cases}
\]
\[
\delta_\beta(i, i+2) = \begin{cases} 
\beta_i, & i = 1, n - 2 \\
\beta_i + \beta_{i+1}, & \text{otherwise}
\end{cases}
\]
\[
\delta_\beta(1, n) = \delta_\beta(2, n) = \beta_2 + \cdots + \beta_{n-2}
\]

We see that \(\delta_\beta(i, i+1) + \delta_\beta(i+1, i+2) = \delta_\beta(i, i+2)\).

Finally, for \(\delta_\gamma\) we further decompose it as \(\delta_\gamma = \delta_{\gamma_1} + \delta_{\gamma_2}\) according to the two summands in (8).

Repeating the same procedure yields
\[
\delta_{\gamma_1}(i, i+1) = \sum_{t=i+2}^{n-1} \gamma_{i,t}
\]
\[
\delta_{\gamma_1}(i, i+2) = \sum_{a=0}^{1} \sum_{t=i+2+a}^{n-1} \gamma_{i+a,t}
\]
\[
\delta_{\gamma_1}(1, n) = \delta_{\gamma_2}(2, n) = 0
\]
\[
\delta_{\gamma_2}(i, i+1) = \delta_{\gamma_2}(i, i+2) = \delta_{\gamma_2}(1, n) = 0
\]
\[
\delta_{\gamma_2}(2, n) = \sum_{t=3}^{n-1} \gamma_{1,t}
\]

After some algebraic manipulations we again obtain \(\delta_\gamma(i, i+1) + \delta_\gamma(i+1, i+2) = \delta_\gamma(i, i+2)\).

**Corollary 9.** Let \(\delta : X \times X \to \mathbb{R}\) be a metric over the finite set \(X\). Then \(\delta\) is circular decomposable if and only if it satisfies the Kalmanson conditions.
4 $\mathcal{K}_n$ and the consecutive ones property

Thus far we have defined $\mathcal{K}_n$ as a split-theoretic simplicial complex and also geometrically in terms of permutations of the Kalmanson conditions. In this section we present a third description of the Kalmanson complex as a set of (equivalence classes of) binary matrices possessing a certain structure. Again, we will show that this formulation is entirely equivalent to the preceding two. Throughout this section, $M$ is taken to be an $m \times n$ binary matrix (entries are zero or one.)

**Definition 5.** $M$ is said to possess the *consecutive ones property for rows* (C1R) if its columns may be permuted such that the ones in each row occur in blocks. $M$ possesses the *circular ones property for rows* (Circ1R) if its columns may be permuted such that either the ones or the zeros (or both) in each row occur in a block.

Intuitively, a Circ1R matrix has the property that for each of its rows, the ones occur in a block when it is “wrapped around” a cylinder.

**Example 3.** Consider the matrices

$$(1) \begin{pmatrix} 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 \end{pmatrix}$$

$$(2) \begin{pmatrix} 0 & 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 & 1 \end{pmatrix}$$

$$(3) \begin{pmatrix} 1 & 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

(1) and (2) are C1R, and (3) is Circ1R. To verify that (2) is C1R, we apply the permutation $(1 \ 3 \ 4 \ 5) \in S_5$ to its columns:

$$(\begin{pmatrix} 0 & 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 & 1 \end{pmatrix})^{(1\ 3\ 4\ 5)} \rightarrow \begin{pmatrix} 0 & 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \end{pmatrix}$$

If $M$ is C1R or Circ1R, then the matrix obtained by replacing any number of rows of $M$ by their binary complement will be Circ1R. This provides justification for the following theorem.
Theorem 10 ([17]). Let $M$ be a binary matrix, and let $M'$ be the matrix obtained by complementing each row in $M$ which has a one in the first column. Then $M$ is C1R if and only if $M'$ is Circ1R.

A circular split system and a Circ1R binary matrix are, in a sense, identical. To see this, let $m$ be fixed and consider the set of all split systems over $X$ which contain $m$ splits: $S_m(X) = \{S \subset S(X) : |S| = m\}$. Also, let $\mathcal{M}_{m \times n}^0(\{0,1\})$ be the set of $m \times n$ binary matrices who first column contains all zeros, and let the symmetric group $S_m$ act on it by permutation of rows.

Finally, let $Q_m = \mathcal{M}_{m \times n}^0(\{0,1\})/\sim$ be the set of equivalence classes under the relation “$M_1 \sim M_2 \iff M_1 = \sigma \cdot M_2$ for some $\sigma \in S_m$”. (Note that, as row permutations do nothing to affect the C1R/Circ1R properties, it makes sense to say that a class $[M] \in Q$ possesses one or both.)

Now define a map $F: S_m(X) \to Q_m$ which sends a system of $m$ splits to the class of the binary matrix obtained by converting the splits to a binary vector and stacking them. Formally,

$$F: S_m(X) \to Q_m$$

$$\left\{ \{A_1,B_1\}, \ldots, \{A_m,B_m\} \right\} \mapsto \left( (w_{ij})_{i \in [m], j \in [n]} \right)$$

$$w_{ij} := \begin{cases} 1, & j \in A_i \text{ and } 1 \notin A_i \\ 1, & j \notin A_i \text{ and } 1 \in A_i \\ 0, & \text{otherwise} \end{cases}$$

Example 4. Let $n = 5$ and $S \in S_3(X)$ be the split system

$$S = \left\{ \left\{ \{1,2\}, \{3,4,5\} \right\}, \left\{ \{1,3,5\}, \{2,4\} \right\}, \left\{ \{1,4\}, \{2,3,5\} \right\} \right\}$$

Then

$$F(S) = \begin{pmatrix} 0 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 1 \end{pmatrix} = \begin{pmatrix} 0 & 1 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 \end{pmatrix}$$

Lemma 11. $F: S_m(X) \to Q_m$ is a bijection.

Proof. Let $[M] \in Q_m$ be given. Simply convert each row of $M \in \mathcal{M}_{m \times n}^0(\{0,1\})$ to an $X$-split in the obvious way. The resulting split system $S$ gives $F(S) = [M]$, so $F$ is onto.

Now suppose $F(S_1) = F(S_2)$ for two split systems $S_1, S_2$. Then $S_1$ and $S_2$ represent the same splits up to ordering. But then $S_1 = S_2$ so $F$ is one-to-one.
Having established the bijection, it is easy to see that $\text{Circ1R}$ and circularity are analogous properties for $Q$ and $\mathcal{S}_m(X)$, respectively.

**Theorem 12.** Let $S \in \mathcal{S}_m(X)$ be an arbitrary split system. Then $S$ is circular iff $F(S)$ is $\text{Circ1R}$.

**Proof.** A split system $S$ is circular iff there is a $\sigma \in S_n$ such that each split $S \in S$ is of the form

$$S = \left\{ \{\sigma(i), \sigma(i+1), \ldots, \sigma(j)\}, \{\sigma(j+1), \ldots, \sigma(i-1)\} \right\}$$

where $\overline{i}$ denotes $i \pmod{n}$. This occurs iff applying $\sigma$ to the columns of $F(S)$ yields a class of matrices $[M] \in Q_m$ whose ones appear consecutively. Since the first column of $M$ is the zero vector, $M$ is $\text{C1P}$ iff it is $\text{Circ1R}$ by Theorem 10.

**Corollary 13.** $S \in \mathcal{S}_m(X)$ is circular iff $F(S)$ is $\text{C1R}$.

The preceding theorem enables us to furnish another description of $\mathcal{K}_n$: it is the poset of all $\text{Circ1R}$ binary matrices (up to row permutation) which possess an initial column of zeros, and at least two ones and two zeros in each row, ordered by inclusion of the set of row vectors corresponding to each matrix.

### 5 $f$-vector

In this section we will harness the three descriptions of $\mathcal{K}_n$ to study its combinatorial structure in greater detail.

**Definition 6.** Let $\Delta$ be a simplicial complex of dimension $d-1$, and let $f_i$ denote the number of $i$-dimensional faces of $\Delta$. The $f$–vector of $\Delta$ is the vector $f = (f_0, f_1, \ldots, f_{d-1})$.

Thus, $f_0$ counts the vertices of $\Delta$ and $f_{d-1}$ counts the facets. By geometric analogy, $f_1$, $f_2$ and $f_{d-2}$ are called the edges, triangles, and ridges of $\Delta$, respectively.

For small $n$, the $f$-vector may be computed directly. Results for $n = 4, \ldots, 9$ are presented in Table 1. We now theoretically explain some of these numbers. First we restate some additional definitions
Table 1: Computational results for the Kalmanson complex.

| $n$ | $f$-vector |
|-----|------------|
| 4   | 3, 3       |
| 5   | 10, 45, 90, 60, 12 |
| 6   | 25, 300, 1755, 4725, 6390, 4860, 2160, 540, 60 |
| 7   | 56, 1540, 19950, 121485, ..., 5040, 360 |
| 8   | 119, 7021, 178878, ..., 50400, 2520 |
| 9   | 246, 30135, 1409590, ..., 544320, 20160 |

$\langle 2^{n-1} - n - 1, \frac{n(n-3)}{4}, \frac{(n-1)!}{2} \rangle$

**Figure 2:** A system of splits which is not weakly compatible.

and results from [4, 10] which will prove useful in enumerating the faces of $K_n$. For the remainder of the section, $S_i \in \mathcal{S}(X)$ represents a split and the identity $S_i = \{A_i, B_i\}$ is implicit.

**Definition 7.** A split system $\mathcal{S}$ is called *weakly compatible* if for all triples $S_1, S_2, S_3 \in \mathcal{S}$ there do not exist points $a, a_1, a_2, a_3 \in X$ such that $a \in A_1 \cap A_2 \cap A_3$ and $a_i \in A_j \iff i = j$.

Weak compatibility enforces a sort of convexity condition on $\mathcal{S}$ by requiring that, for any triple of points in the split system, there is no point which mutually separates them (Figure 2).

For any two splits $S_1, S_2$ we define a binary operation $\sqcup$ by $\{A_1, B_1\} \sqcup \{A_2, B_2\} = \{A_1 \cap A_2, B_1 \cup B_2\}$.
Lemma 14. The splits $S_1$, $S_2$ and $S_1 \sqcup S_2$ are weakly compatible.

Proof. Let $S_3 = S_1 \sqcup S_2$. If there exist $a,a_1,a_2,a_3$ as in Definition 7, then $a_3 \in A_3 - (A_1 \cup A_2) \neq \emptyset$. But $A_3 = A_1 \cap A_2$, a contradiction. \hfill \Box

Theorem 15 ([10]). Let $S$ be a split system and let $S'$ be the split system

$$S' := S \cup \left\{ S_1 \sqcup S_2 : S_1, S_2 \in S \text{ and } A_i \cap B_j \neq 0, \ i, j \in \{1, 2\} \right\}$$

Then $S$ is contained in a circular split system if and only if $S'$ is weakly compatible.

Corollary 16. A circular split system is weakly compatible.

5.1 Low (Co-)Dimensional Faces

Enumerating the vertices, edges, ridges and facets of $K_n$ is now straightforward.

Theorem 17. Let $f = (f_0, \ldots, f_{d-1})$ denote the $f$-vector of $K_n$. Then

$$f_0 = 2^{n-1} - n - 1 \quad (9)$$
$$f_1 = \binom{f_0}{2} \quad (10)$$
$$f_{d-2} = \left[ \binom{n}{2} - n \right] \times f_{d-1} \quad (11)$$
$$f_{d-1} = \frac{(n-1)!}{2} \quad (12)$$

Proof. $f_0$ counts the number of non-trivial $X$-splits. There are

$$\sum_{k=2}^{n-2} \binom{n}{k} = 2^n - 2n - 2$$

binary words on $n$ letters which contain at least two zeros and two ones. Since each word and its complement correspond to the same split, we divide by two to obtain $f_0$.

Equation (10) asserts that every pair of splits $S_1, S_2$ is contained in a circular split system. By Lemma 14, the splits $S_1$, $S_2$ and $S_1 \sqcup S_2$ are weakly compatible. Then by Theorem 15, $\{S_1, S_2\}$ is a circular split system.
Each facet of $K_n$ corresponds to a circular ordering; that is, an edge labeling of the regular $n$-gon. Such labelings are unique up to dihedral symmetry. There are $(n - 1)!$ labelings up to rotation, and half that number when accounting for reflection. This yields (12).

To prove (11), let $F \subset K_n$ be a facet spanned by vertices $v_1, \ldots, v_d \in X$; without loss of generality assume the circular ordering corresponding to $F$ is $(1, 2, \ldots, n)$. Let $u$ be another vertex distinct from the $v_i$, with corresponding split

$$S_u = \left\{ \{1, u_2, \ldots, u_j\}, \{u_{j+1}, \ldots, u_n\} \right\}$$

Finally, let $i = \min \{i \colon u_i \neq i\}$, which exists by the assumption that $S_u$ is not circular with respect to the given ordering. Now, the splits $S_u$ and

$$S_1 = \left\{ \{u_i - 1, u_i\}, X - \{u_i - 1, u_i\} \right\} \in F$$

$$S_2 = \left\{ \{u_i, u_i + 1\}, X - \{u_i, u_i + 1\} \right\} \in F$$

are weakly incompatible: denoting the first blocks of each by $A_u, A_1, A_2$ we have

$$\{u_i\} = A_u \cap A_1 \cap A_2$$

$$1 \in A_u - (A_1 \cup A_2)$$

$$u_i - 1 \in A_1 - (A_u \cup A_2)$$

$$u_i + 1 \in A_2 - (A_u \cup A_1)$$

Hence, by contradiction any collection of $d - 1$ vertices of $F$ spans a unique face of codimension two. As described in Section 3 each facet contains $\binom{n}{2} - n$ vertices (one for each diagonal of the $n$-gon.)

5.2 Triangles

The computations in Theorem 17 were aided by the fact that $K_n$ is connected in dimension one and totally disconnected in codimension one. Enumerating the faces in the remaining cases is more challenging. To illustrate the issues involved, we demonstrate how to compute $f_2$, the number of triangles in $K_n$.

**Example 5.** The split system

$$\left\{ \left\{ \{1, 2\}, \{3, 4, 5\} \right\}, \left\{ \{1, 3\}, \{2, 4, 5\} \right\}, \left\{ \{1, 4\}, \{2, 3, 5\} \right\} \right\}$$
is not weakly compatible, so it is not a triangle of $K_n$. By contrast, the split system

$$\left\{ \left\{ \{1, 2\}, \{3, 4, 5\} \right\}, \left\{ \{2, 3\}, \{1, 4, 5\} \right\}, \left\{ \{4, 5\}, \{1, 2, 3\} \right\} \right\}$$

is circular with respect to two orderings: $(1, 2, 3, 4, 5)$ and $(1, 2, 3, 5, 4)$. It is therefore a triangle of $K_n$ which is contained in two facets.

Our main tool for computing $f_2$ will be Corollary 13, in conjunction with a structure theorem of [18] which completely characterizes C1R matrices.

**Definition 8.** Let $M$ be a matrix. The *configuration of* $M$ is the set of matrices obtained by permuting the rows and/or columns of $M$ (not necessarily by the same permutation.)

**Example 6.** The configuration of the $2 \times 2$ identity matrix is the set

$$\left\{ \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \right\}$$

**Theorem 18 ([18]).** A binary matrix $M$ is C1R if and only if it does not contain as a submatrix any configuration of $M_{I_n}, M_{II_n}, M_{III_n}, M_{IV}, M_V, 1 \leq n < \infty$, where

$$M_{I_n} = \begin{pmatrix} r_1 & \begin{pmatrix} c_1 & c_2 & c_3 & \cdots & c_n & c_{n+1} & c_{n+2} \\ 1 & 1 & 0 & \cdots & 0 & 0 & 0 \\ 0 & 1 & 1 & \cdots & 0 & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \cdots & 1 & 1 & 0 \\ 0 & 0 & 0 & \cdots & 0 & 1 & 1 \\ 1 & 0 & 0 & \cdots & 0 & 0 & 1 \end{pmatrix} & r_2 \\ r_3 \\ r_4 \end{pmatrix} \quad M_{IV} = \begin{pmatrix} r_1 & \begin{pmatrix} c_1 & c_2 & c_3 & c_4 & c_5 & c_6 \\ 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 \end{pmatrix} & r_2 \\ r_3 \\ r_4 \end{pmatrix}$$

$$M_{II_n} = \begin{pmatrix} r_1 & \begin{pmatrix} c_1 & c_2 & c_3 & \cdots & c_n & c_{n+1} & c_{n+2} & c_{n+3} \\ 1 & 1 & 0 & \cdots & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & \cdots & 0 & 0 & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \cdots & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & \cdots & 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & \cdots & 1 & 1 & 0 & 1 \\ 0 & 1 & 1 & \cdots & 1 & 1 & 1 & 1 \end{pmatrix} & r_2 \\ r_3 \\ r_4 \end{pmatrix} \quad M_{V} = \begin{pmatrix} r_1 & \begin{pmatrix} c_1 & c_2 & c_3 & c_4 & c_5 \\ 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 & 1 \end{pmatrix} & r_2 \\ r_3 \\ r_4 \end{pmatrix}$$

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In the specific case of $f_2$, where we are counting $3 \times n$ matrices, only two forbidden submatrices pertain:

$$M_{I_1} = \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 1 \end{pmatrix} \quad \text{and} \quad M_{III_1} = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 \end{pmatrix}$$

For a matrix $M$, we write $\text{col}(M)$ to denote the set of column vectors of $M$. Let $I = \text{col}(M_{I_1})$ and $III = \text{col}(M_{III_1})$. Note that $I$ and $III$ are “closed” under the operation of row permutation. Hence, by Corollary 13 and Theorem 18,

$$[M] \in \mathcal{Q}_3 \iff |\text{col}(M) \cap I| < 3 \quad \text{and} \quad |\text{col}(M) \cap III| < 4$$

Accordingly, let

$$F_{i,j} = \{[M] \in \mathcal{Q}_3 : |\text{col}(M) \cap I| = i \quad \text{and} \quad |\text{col}(M) \cap III| = j\}$$

Then

$$f_2 = |\mathcal{Q}_3| = \sum_{0 \leq i \leq 2} \sum_{0 \leq j \leq 3} |F_{i,j}|$$

Enumerating $F_{i,j}$ involves carefully counting the number of classes of $\mathcal{Q}_3$ while keeping track of how many columns from the sets

$$I = \left\{ \begin{pmatrix} 1 \\ 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \\ 1 \end{pmatrix} \right\} \quad \text{and} \quad III = \left\{ \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} \right\}$$

appear in each equivalence class.
5.2.1 Sample Calculation: $|F_{0,3}|$

The counting argument is straightforward but tedious. We illustrate the calculation of $|F_{0,3}|$; the remaining cases are similar and are proven in [16].

Let $\mathbb{P}_n$ denote the set of ordered partitions of the integer $n$. That is, for a $k$-tuple $(x_1, \ldots, x_k)$ we have

$$(x_1, \ldots, x_k) \in \mathbb{P}_n \iff \sum_{i=1}^{k} x_i = n \text{ and } x_i \geq 1 \text{ for all } i$$

To simplify the notation we take summation over $\mathbb{P}_{n-1}$ for granted wherever there is no chance of confusion: instead of e.g.

$$\sum_{(a,b,c,d) \in \mathbb{P}_{n-1}} \binom{n-1}{a,b,c,d}$$

we will simply write

$$\sum_{a>1} \binom{n-1}{a,b,c,d}$$

Now let $[M] \in F_{0,3}$. We consider two cases.

1. First, if $(1,1,1)^T \notin \text{col}(M) \cap III$ then $M$ is of the form

$$
\begin{pmatrix}
  a & b & c \\
  0 & 0 & 0 \\
  0 & 0 & 0 \\
  0 & 1 & 1 \\
\end{pmatrix}
\begin{pmatrix}
  1 & 1 & 0 \\
  0 & 0 & 0 \\
  0 & 1 & 1 \\
\end{pmatrix}
$$

where $a, b, c$ count the instances of the columns in $III - \{(1,1,1)\}$. To prevent the occurrence of a trivial split (row containing $< 2$ ones) we require $\min(a,b,c) > 1$. Hence there are

$$(1/6) \sum_{\min(a,b,c)>1} \binom{n-1}{a,b,c}$$

such classes, where the factor of $1/6$ reflects the fact that each arrangement is equivalent to six others obtained by permuting the labelings $a, b, c$. We see that this counts the number ways of arranging $n - 1$ ones into three unlabeled rows where each must contain at least two ones.\footnote{This number is also known as an associated Stirling number of the second kind, cf. A000478 [13].}
We also have the possibility that $M$ contains additional zero columns:

\[
\begin{pmatrix}
  d & a & b & c \\
  0 & 0 & 1 & 1 & 0 & 0 & 0 & 0 \\
  0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 \\
  0 & 0 & 0 & 0 & 0 & 0 & 1 & 1
\end{pmatrix}
\]

By an entirely analogous argument we count

\[
\frac{1}{6} \sum_{\min(a,b,c)>1} \binom{n-1}{a,b,c,d}
\]

such classes.

2. In the second case we have $(1,1,1)^T \in \text{col}(M) \cap III$. These are matrices of the form

\[
\begin{pmatrix}
  c & a & b \\
  0 & 1 & 1 & 1 & 0 & 0 & 1 & 1 \\
  0 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\
  0 & 1 & 1 & 1 & 1 & 1 & 0 & 0
\end{pmatrix}
\]

We must have that the number of $(1,1,1)^T$ columns is greater than one, or else we have a trivial split. Also, each arrangement is equivalent to the one obtained by swapping the columns labeled $a$ and $b$ and permuting their respective rows. We therefore count

\[
\frac{1}{2} \sum_{c>1} \binom{n-1}{a,b,c}
\]

such classes. We also obtain

\[
\frac{1}{2} \sum_{c>1} \binom{n-1}{a,b,c,d}
\]

classes by allowing for the presence of additional zero columns.

Now let $S_{n,k}$ denote the Stirling number of the second kind. We will make use of the identity

\[
\sum_{(x_1,\ldots,x_k)\in\mathbb{P}_n} \binom{n}{x_1,\ldots,x_k} = k! \cdot S_{n,k} =: M(n,k),
\]

(19)

to obtain simple formulas for equations (15)–(18). (One interpretation of $M(n,k)$ is that it counts the number of surjections from an $n$-set onto a $k$-set, see e.g. [1, Ch. 1].) By (19), symmetry and

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the inclusion-exclusion principle we have

\[
\sum_{\min(a,b,c) > 1} \binom{n-1}{a,b,c} = M(n-1,3) - 3 \sum_{a=1}^{n-1} \binom{n-1}{a,b,c} + 3 \sum_{a=b=1}^{n-1} \binom{n-1}{a,b,c} - \sum_{a=b=c=1}^{n-1} \binom{n-1}{a,b,c}
\] (20)

The second term of (20) counts the number of surjections from a set of cardinality \(n - 1\) onto the set \(\{a, b, c\}\) such that a unique element maps to \(a\). There are \((n - 1) \times M(n - 2, 2)\) such maps. Similarly, the third and fourth terms represent \(2! \times \binom{n-1}{2} \times M(n - 3, 1) = (n - 1)(n - 2)\) and \(3! \times (n - 1, 3) \times M(n - 4, 0) = 0\) maps respectively. We therefore conclude

\[
\sum_{\min(a,b,c) > 1} \binom{n-1}{a,b,c} = M(n-1,3) - 3(n-1)M(n-2,2) + 3(n-1)(n-2)
\]

By the same arguments,

\[
\sum_{\min(a,b,c,d) > 1} \binom{n-1}{a,b,c,d} = M(n-1,4) - 3(n-1)M(n-2,3) + 3(n-1)(n-2)M(n-3,2) - (n - 1)(n - 2)(n - 3)
\]

For (17) we note that

\[
\sum_{c > 1} \binom{n-1}{a,b,c} = M(n-1,3) - \sum_{c=1}^{n-1} \binom{n-1}{a,b,c} = M(n-1,3) - (n - 1) \times M(n - 2, 2)
\]

Similarly, we can rewrite (18) as \(M(n-1,4) - (n - 1) \times M(n - 2, 3)\).

5.2.2 General Formula

Repeating these counting arguments for the remaining \(|F_{i,j}|\) yields the following formula.

**Theorem 19.** Let \(t = n - 1\). The number of triangles in \(K_n\) is

\[
(1/6)(t-2)(t-1)t + 2(t-1)t [1 + M(t - 2, 2)] - 5t \cdot M(t - 1, 2) - 8t \cdot M(t - 1, 3) - 2t \cdot M(t - 1, 4) + (19/6) M(t, 3) + (55/6) M(t, 4) + 7M(t, 5) + 2M(t, 6)
\]

**Proof.** See [16].
The first ten entries of this sequence, $n = 4, \ldots, 13$ are

$$0, 90, 1755, 19950, 178878, 1409590, 10270585, 71110930, 475443364, 3100707610, \ldots$$

We have verified this formula computationally up to $n = 10$ (the largest $n$ for which the calculations terminated) using the mathematics software SAGE [15]. Source code for this and related $f$-vector calculations may be downloaded from: https://github.com/terhorst/kalmanson

6 Conclusion

In this paper we analyzed the combinatorics of the Kalmanson complex. We show how this complex arises in split theory, optimization and phylogenetics. We gave a simplified proof of the equivalence of Kalmanson and circular decomposable metrics based on polyhedral geometry and our interpretation of $K_n$ as a simplicial complex of splits.

Subsequently, our main focus was to enumerate its faces. We demonstrated that the complex is totally connected along edges, and totally disconnected along ridges. We then gave a formula for enumerating its triangles, using a forbidden substructure characterization along with some basic counting principles.

At present we do not have a way to generalize this method to faces of arbitrary dimension. The next case of tetrahedra ($k = 4$) becomes considerably more difficult, as there are now 7 avoided Tucker matrices to consider: $M_{I_1}, M_{I_2}, M_{II_1}, M_{III_1}, M_{III_2}, M_{IV}, M_{V}$. The connection to the Tucker theorem suggests a possible application of results on avoided configurations (see [2] for a survey), but most results in that literature are of an extremal, as opposed to enumerative, variety. In [14] some matrices avoiding small configurations are counted, but we are not aware of a general method of enumerating matrices which avoid configurations of arbitrary dimensions. We view this as an interesting problem in enumerative combinatorics which merits further study.
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