AN ALGORITHM FOR PRODUCING F-PURE IDEALS

ALBERTO F. BOIX* AND MORDECHAI KATZMAN**

Abstract. This paper describes a method for computing all F-pure ideals for a given Cartier map of a polynomial ring over a finite field.

Introduction

The subject of this paper is the study of certain ideals associated with a given $p^{-\epsilon}$-linear map. These maps were introduced by K. Schwede in [11] and M. Blickle in [1] in the context of test ideals and are defined as follows.

Throughout this manuscript, unless otherwise is specified, we shall denote by $A$ a fixed regular ring which is either the polynomial ring $\mathbb{K}[x_1, \ldots, x_d]$, the localization $\mathbb{K}[x_1, \ldots, x_d]_m$ (where $m$ is the ideal generated by all the variables), or the formal power series ring $\mathbb{K}[[x_1, \ldots, x_d]]$, where $\mathbb{K}$ is an $F$-finite field of prime characteristic $p$ (i.e. a field $\mathbb{K}$ which is a finite extension of $\mathbb{K}^p$). The Frobenius map, raising an element $a \in A$ to its $p$th power $a^p$, is an additive map.

Given any $A$-module $M$ and $\epsilon \in \mathbb{N}$, $F^\epsilon_* M$ shall denote the abelian group $M$ with $A$-module structure given by $a \cdot m = a^p m$ for any $a \in A$ and $m \in F^\epsilon_* M$. Given an $m \in M$ we shall henceforth write $F^\epsilon_* m$ for the same element regarded as a member of $F^\epsilon_* M$.

The $p^{-\epsilon}$-linear maps referred to above are elements in $\text{Hom}_A(F^\epsilon_* M, M)$; these can be thought as additive maps $M \xrightarrow{\phi_\epsilon} M$ for which, for any $a \in A$ and $m \in M$, $\phi_\epsilon(a^p m) = a\phi_\epsilon(m)$. We further define

$$C^M := \bigoplus_{\epsilon \geq 0} \text{Hom}_A(F^\epsilon_* M, M)$$

and endow it with the structure of an $A$-algebra by defining the product of $\phi_\epsilon \in C^M$ and $\phi_\epsilon' \in C^M$ as the element of $C^M$ given by

$$F^\epsilon + \epsilon' \xrightarrow{\phi_\epsilon + \epsilon'} M.$$

Moreover, we also set

$$C^M_+ := \bigoplus_{\epsilon \geq 1} \text{Hom}_A(F^\epsilon_* M, M).$$

In this article we shall be interested mostly in $C^A$.

Definition. A $C^M$-submodule $N$ of $M$ is $F$-pure if $C^M_+ N = N$. In particular, we say that an ideal $I$ of $A$ is $F$-pure provided $C^A_+ I = I$.

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Our paper is motivated by the study of $F$-pure ideals and their properties introduced in [1]. When $A$ is $F$-finite, $A$ itself is an $F$-pure ideal if and only if $C_A^+$ contains a splitting of a certain power of the Frobenius map on $A$ (cf. [1 Proposition 3.5]); therefore, these $F$-pure ideals turn out to be a generalization of the $F$-purity property.

Furthermore, among the main results in [1] (see also [2, Corollary 4.20 and Proposition 5.4]) is the fact that the set of $F$-pure ideals in $A$ is finite, and that the big test ideal is the minimal element of a set of $F$-pure ideals. Apart from their usefulness in describing big test ideals, we believe that the set of $F$-pure ideals provides an interesting set of invariants of $A$ providing information about the ring which is not yet fully understood.

One should contrast this with the situation one encounters where studying the set of $C_A^+$-compatible ideals, i.e., ideals $I \subseteq A$ for which $C_A^+I \subseteq I$. One might hope to list all $F$-pure ideals by listing all compatible ideals and checking which ones are $F$-pure. However, the set of compatible ideals need not be finite, and one can only describe algorithmically the radical ideals among these; such task was carried out in [9].

Our contribution to the understanding of $F$-pure ideals is to provide an effective procedure to calculate all the $F$-pure ideals contained in $m$ of the subalgebra $C = C^\phi$ of $C_A^\phi$ generated by one homogeneous element $\phi \in \text{Hom}_A(F^*_e A, A)$ under the additional assumption that the ground field $K$ is finite. This procedure has been implemented in Macaulay2 (cf. [4]).

Now, we provide an overview of the contents of this manuscript. Firstly, in Section 1 we introduce compatible and fixed ideals; moreover, we show that the so-called $\epsilon$th root ideal (cf. Definition 1.2) plays a key role in their calculation. Secondly, Section 2 contains the main result of this paper; namely, the algorithm referred to above (cf. Theorem 2.5). By the way, we introduce a new operation on ideals (cf. Definition 2.2), hoping that it may be interesting in its own right. Finally, in Section 3 we provide examples in order to illustrate how our method works; most of these specific computations were carried out with an implementation of this procedure in Macaulay2.

1. Ideals compatible and fixed under a given $p^{-\epsilon}$-linear map

Unless otherwise is specified, $K$ is to denote an $F$-finite field of prime characteristic $p$, i.e., a field $K$ which is a finite extension of $K^p$. Given an $\alpha = (a_1, \ldots, a_d) \in \mathbb{N}^d$ we shall use the following multi-index notation:

$$\mathbf{x}^\alpha := x_1^{a_1} \cdots x_d^{a_d}.$$  

Moreover, in this case, we set $||\mathbf{x}^\alpha|| := \max\{a_1, \ldots, a_d\}$ and, for any polynomial $g \in K[x_1, \ldots, x_d]$, 

$$||g|| := \max_{\alpha \in \text{supp}(g)} ||\mathbf{x}^\alpha||,$$

where $g = \sum_{\alpha \in \mathbb{N}^d} g_{\alpha} \mathbf{x}^\alpha$ (such that $g_{\alpha} = 0$ up to a finite number of terms) and 

$$\text{supp}(g) := \{\alpha \in \mathbb{N}^d \mid g_{\alpha} \neq 0\}.$$  

Finally, given any ideal $I$, $I[p^\epsilon]$ shall denote the ideal generated by all $p^\epsilon$ powers of elements in $I$. It is straightforward to verify that $I[p^\epsilon]$ is generated by the $p^\epsilon$ powers of a set of generators of $I$. 
The $F$-finiteness of $K$ implies that $F^r_e A$ is a free $A$-module of finite rank. Indeed, if $B_e$ is a $K^{p^r}$-basis for $K$, then $F^r_e A$ has free basis

$$\{b \mathbf{x}^\alpha \mid b \in B_e, \quad 0 \leq ||\alpha|| \leq p^r - 1\}.$$ 

Moreover, we recall that the trace map $\Phi_e \in \text{Hom}_A(F^r_e A, A)$, which is the projection onto the direct summand $Ax_1^{p^r-1} \cdots x_d^{p^r-1}$, generates the $F^r_e A$-module $\text{Hom}_A(F^r_e A, A)$ (cf. [3, Example 1.3.1]). In this way, any homogeneous element $\phi \in \text{Hom}_A(F^r_e A, A)$ can be written as $u\Phi_e$ (interpreted as the composition of multiplication by $u$ followed by $\Phi_e$) for some $u \in F^r_e A$. Now if $C$ is the Cartier subalgebra of $C^A$ generated by such a $\phi$ then the problem of finding the $F$-pure ideals of $A$ amounts to finding all ideals $I \subseteq A$ such that $\phi(F^r_e I) = I$.

**Definition 1.1.** Let $I$ be an ideal of $A$ and let $\phi \in \text{Hom}_A(F^r_e A, A)$.

(i) We say that $I$ is $\phi$-compatible if $\phi(F^r_e I) \subseteq I$.

(ii) We say that $I$ is $\phi$-fixed if $\phi(F^r_e I) = I$.

Clearly, all $\phi$-fixed ideals are $\phi$-compatible. The converse also holds if $\phi$ is a Frobenius splitting, i.e., if $\phi(F^r_e 1) = 1$: in this case, for any $r$ in a $\phi$-compatible ideal $I$ we have $\phi(F^r_e p^r) = r \phi(F^r_e 1) = r$.

From now on, we shall write our given $\phi \in \text{Hom}_A(F^r_e A, A)$ as $u\Phi_e$, where $u \in F^r_e A$.

1.1. **The ideal of $p^r$-th roots.** Our next goal is to express in an equivalent way the condition of being $\phi$-fixed in order to perform explicit calculations. Such an equivalent expression requires us to introduce the following concept (cf. [3, Definition 2.2] and [7, Section 5]).

**Definition 1.2.** Let $J$ be an ideal of $A$. We set $I_e(J)$ as the smallest ideal $I$ such that $I^{[p^r]} \supseteq J$. We shall refer to $I_e(J)$ as the $e$-th root ideal of $J$ (it is sometimes denoted $J^{[1/p^r]}$).

We have the following elementary properties of $e$-th roots (see [7, Section 5] for details).

**Proposition 1.3.** Let $J, J_1, \ldots, J_r$ be ideals of $A$. Then, the following statements hold.

(a) If $J_1 \subseteq J_2$ then $I_e(J_1) \subseteq I_e(J_2)$.

(b) One has that

$$I_e \left( \sum_{i=1}^r J_i \right) = \sum_{i=1}^r I_e(J_i).$$

Note that this fact implies that it is enough to know how to calculate $I_e(J)$ when $J$ is a principal ideal.

(c) Let $g \in A$. If

$$g = \sum_{\substack{b \in B_e \\alpha \in \mathbb{Z}^d \\alpha \leq ||\alpha|| \leq p^r - 1}} g_{ab}^p b \mathbf{x}^\alpha$$

then $I_e(g)$ is the ideal of $A$ generated by all the $g_{ab}$’s.

Now, we are ready for expressing the condition of being $\phi$-fixed in computational terms. This is the main result of this section.
Theorem 1.4. Let $J \subseteq A$ be any ideal and let $\phi = u\Phi_e \in \text{Hom}_A(F_e^* A, A)$. Then, the following statements hold.

(a) The image of $F_e^* J$ under $\phi$ is $I_e(uJ)$.
(b) $J$ is $\phi$-compatible if and only if $I_e(uJ) \subseteq J$.
(c) $J$ is $\phi$-fixed if and only if $I_e(uJ) = J$.

Proof. Parts (b) and (c) follow directly from part (a). So, it is enough to prove part (a).

Proposition 1.3 implies that, in order to compute $I_e(uJ)$, one may choose a set of generators $g_1, \ldots, g_t$ of $F_e^* J$ and then compute $I_e(ug_1) + \ldots + I_e(ug_t)$. Now, fix $1 \leq i \leq t$ and write

$$ug_i = \sum_{b \in \mathcal{B}_e} r_{iab}^e b^\alpha x^\alpha.$$ 

Applying once more Proposition 1.3 it follows that $I_e(ug_i)$ is the ideal generated by all coefficients $r_{iab}$ above. But $r_{iab} = \Phi_e(F_e^*(b^e x_1^{p^e-\alpha_1} \ldots x_d^{p^e-\alpha_d}) u g_i) \in \phi(F_e^* J)$, hence $I_e(ug_i) \subseteq \phi(F_e^* J)$ for any $1 \leq i \leq t$ and $I_e(uJ) \subseteq \phi(F_e^* J)$.

Conversely, note that $\phi(y) = \Phi_e(uy) \in I_e(uJ)$ for any $y \in F_e^* J$, hence $\phi(F_e^* J) \subseteq I_e(uJ)$ and therefore we obtain the desired conclusion. \(\square\)

Notation 1.5. Henceforth $S$ shall denote the polynomial ring $K[x_1, \ldots, x_d]$ and $S_l$ shall denote the $K$-vector space generated by monomials $x^\alpha$ with $||\alpha|| \leq l$.

The following result will guarantee that the algorithm we shall introduce later on terminates after a finite number of steps.

Proposition 1.6. The following statements hold.

(i) For any $y \in S$, the ideal $I_e(y)$ can be generated by elements $g \in S$ such that

$$||g|| \leq \frac{||y||}{p^e}.$$ 

(ii) If $J$ is $u\Phi_e$-fixed then there exists a set of generators of $J$ such that if $g$ belongs to such set then

$$||g|| \leq \frac{||u||}{p^e - 1}.$$ 

Proof. We begin proving part (i). Indeed, we write

$$y = \sum_{b \in \mathcal{B}_e} \sum_{0 \leq ||\alpha|| \leq p^e - 1} y_{ab}^e b^\alpha x^\alpha.$$ 

In this way, for any $\alpha$ and $b$ as above it follows that

$$p^e ||y_{ab}|| \leq ||y_{ab}^e b^\alpha|| \leq ||y_{ab}^e x^\alpha|| \leq ||y||,$$

whence part (i) holds.

Now, we prove part (b). Let $M \geq 0$ be the minimal integer for which a set of generators of $J$ have norm at most $M$. Part (i) shows that $I_e(uJ)$ can be generated by polynomials with norm at most $(||u|| + M)/p^e$. In addition, as $I_e(uJ) = J$
we deduce, by the minimality of $M$, that $M \leq (||u|| + M)/p^e$ and therefore we conclude that $M \leq ||u||/(p^e - 1)$, just what we finally wanted to check. \qed

2. The algorithm through the hash operation

The aim of this section is to describe a computational method to produce all the $u\Phi_e$-fixed ideals of $S$. As the reader will appreciate, our procedure is based on a new operation on ideals (cf. Definition 2.2), which we hope to be of some interest in its own right.

We start with the following elementary statement, which we provide a proof for the sake of completeness. It may be regarded as an elementary consequence of Nakayama’s Lemma.

**Lemma 2.1.** Let $I \subseteq m$ be an ideal minimally generated by $s$ elements. Then, any ideal $J \varsubsetneq I$ is contained in some ideal $V$, where $mI \subseteq V \subseteq I$ and $\dim_K I/V = 1$.

**Proof.** Nakayama’s Lemma implies that there are $g_1, \ldots, g_s \in S$ with $I = S g_1 + \ldots + S g_s$ such that $g_1, \ldots, g_s$ (mod $mK$) is a basis of the $s$-dimensional $K$-vector space $I/mI$. In this way, it follows that any ideal $J \subseteq I$ is contained in some $V := SW + mI$, where $W$ is a $(s - 1)$-dimensional $K$-vector subspace of $I/mI$. Moreover, we have to note as well that $\dim_K I/V = 1$. \qed

From now on, we shall assume that $u \in F^* S$ is fixed and set

$$D_e := \left\lceil \frac{||u||}{p^e - 1} \right\rceil.$$  

The following construction will become in the crucial building block of our method.

**Definition 2.2.** Given any ideal $J \subseteq S$, we define the sequence of ideals

$$J_0 := J, \quad J_{i+1} := \left( J_i \cap \left( J_i^{[p^e]} : S u \right) \cap I_e(uJ_i) \cap S \right) S,$$

and set

$$J^\# := \bigcap_{i \geq 0} J_i.$$  

When $e = 1$, we shall write $J^\#$ instead of $J^\#_1$ for the sake of brevity. Hereafter, we are to refer to this construction as the hash operation.

The introduction of the hash operation is motivated by the following result.

**Lemma 2.3.** For any ideal $J \subseteq S$, $J^\#$ contains all the $u\Phi_e$-fixed ideals which are contained in $J$.

**Proof.** Let $I \subseteq J$ be any $u\Phi_e$-fixed ideal. We shall show by increasing induction on $i \geq 0$ that $I \subseteq J_i$, where $J_i$ is as in Definition 2.2. This is clearly true for $i = 0$.

Now, we assume that $i \geq 0$ and that $I \subseteq J_i$. First of all, as $I$ is $u\Phi_e$-fixed it is, in particular, $u\Phi_e$-compatible and therefore

$$uI \subseteq I^{[p^e]} \subseteq J_i^{[p^e]},$$

hence $I \subseteq \left( J_i^{[p^e]} : S u \right)$. Secondly, using once more that $I$ is $u\Phi_e$-fixed it follows that $I = I_e(uI) \subseteq I_e(uJ_i)$ (indeed, we are using as well that $I \subseteq J_i$ and that the
eth root operation on ideals preserves inclusions). In this way, the foregoing two facts allow us to say that
\[ I \subseteq J_i \cap \left( J_i^{[p^n]} :_{S} u \right) \cap L_e(uJ_i). \]

Finally, bearing in mind that \( I = (I \cap S_{D_e})S \) (indeed, here we are applying part (ii) of Proposition 2.6 and the previous facts it follows that
\[ I = (I \cap S_{D_e})S \subseteq J_i, \]
just what we finally wanted to check. □

**Remark 2.4.** As we have seen in Lemma 2.4, the hash operation produces the smallest compatible ideal \( J^\# \) contained in a given one (namely, \( J \)) such that all the \( \phi \)-fixed ideals contained in \( J \) are also contained in \( J^\# \). This fact allows us to regard this construction as a kind of round down operation on ideals.

### 2.1. The statement of the algorithm.

Now, we introduce our promised algorithm. More precisely, the next result is a recursive procedure for producing all the \( u\Phi_e \)-fixed ideals of \( S \).

This is the main result of this paper.

**Theorem 2.5.** Let \( I \subseteq m \). The set \( \text{FP}_e(I) \) of all \( u\Phi_e \)-fixed ideals contained in \( I \) is given recursively as \( \text{FP}_e(\langle 0 \rangle) = \{ \langle 0 \rangle \} \) and, for \( I \neq \langle 0 \rangle \), defined as the union of \( \{ I^\# \} \) (whenever \( I^\# \) is \( u\Phi_e \)-fixed) and
\[
\bigcup \left\{ \text{FP}_e(V) \mid mI^\# \subseteq V \subseteq I^\#, \quad \dim_K I^\# / V = 1 \right\}.
\]
Moreover, if \( K \) is finite then this recursion is finite in the sense that the resulting execution tree is finite.

**Proof.** Firstly, we show that if \( J \subseteq I \) is \( u\Phi_e \)-fixed then \( J \in \text{FP}_e(I) \). We shall proceed by increasing induction on \( t := \dim_K(I \cap S_{D_e}) \).

Indeed, if \( t = 0 \) then \( J \subseteq I^\# = \langle 0 \rangle \) and therefore \( J \in \{ \langle 0 \rangle \} = \text{FP}_e(I) \).

Now, let \( J \subseteq I \) be such that \( t \geq 1 \). If \( J = I^\# \) then we are done by Lemma 2.4. Thus, we assume that \( J \subseteq I^\# \). Since \( I \subseteq m \), Lemma 2.4 says us that we can find an ideal \( mI^\# \subseteq V \subseteq I^\# \) such that \( \dim_K I^\# / V = 1 \) and \( J \subseteq V \). Furthermore, by construction, \( I^\# \) can be generated by elements in \( S_{D_e} \), hence \( V \cap S_{D_e} \subseteq I^\# \cap S_{D_e} \) and therefore the induction hypothesis implies that \( J \in \text{FP}_e(V) \subseteq \text{FP}_e(I) \).

Finally, we have to point out that our foregoing inductive argument shows that the chains of \( V^\# \)'s produced in this recursion have length at most \( \dim_K S_{D_e} \), hence the second statement follows too. □

In this way, we can turn Theorem 2.5 into an effective method to calculate all the \( u\Phi_e \)-fixed ideals of any polynomial ring having a finite field as field of coefficients as follows.

**Algorithm 2.6.** Let \( K \) be a finite field of prime characteristic, set \( S := K[x_1, \ldots, x_d] \) and let \( u \in S \). These data act as the input of the procedure. Moreover, we initialize \( I \) as the whole ring \( S \) and \( L \) as the empty list \( \{ \} \).

(i) Compute \( I^\# \). Assign to \( I \) the value of \( I^\# \).
(ii) If \( L_e(uI) = I \), then add \( I \) to the list \( L \).
(iii) If \( I = 0 \), then stop and output the list \( L \).
(iv) If \( I \neq 0 \) but principal, assign to \( I \) the value of \( mI \) and loop over the previous steps.
(v) If \( I \neq 0 \) and not principal, then compute
\[
\{ V \text{ ideal} \mid mI \subseteq V \subseteq I, \dim_K I/V = 1 \}.
\]

For each element \( V \) of such set, loop over the previous steps.

At the end of this method, the list \( L \) contains all the \( u\Phi_e \)-fixed ideals of \( S \).

Remark 2.7. The reader should notice that step (v) of the previous method is the only reason for which we have to assume that our coefficient field \( K \) is finite; otherwise, the set \( \{ V \text{ ideal} \mid mI \subseteq V \subseteq I, \dim_K I/V = 1 \} \) is not finite.

We end this section with the following result, which may be regarded as an elementary consequence of the very definition of the hash operation.

**Corollary 2.8.** Let \( K \) be any \( F \)-finite field of prime characteristic \( p \), set \( S := K[x_1, \ldots, x_d] \), and let \( u \in S \). Then, the ideal \( \langle u \rangle \) is a minimal \( \phi \)-fixed ideal, where \( \phi := u^{p-1} \Phi_1 \).

**Proof.** We have to check that \( \langle u \rangle \) is a minimal non-zero \( \phi \)-fixed ideal of \( S \). Firstly, we are to show that \( \langle u \rangle \) is \( \phi \)-fixed; indeed,
\[
I_1(u^{p-1} \cdot \langle u \rangle) = I_1(u^p) = \langle u \rangle,
\]
whence \( \langle u \rangle \) is \( \phi \)-fixed. So, it only remains to prove that \( \langle u \rangle \) is a minimal \( \phi \)-fixed ideal of \( S \). First of all, notice that
\[
D_1 = \left\lceil \frac{\|u^{p-1}\|}{p-1} \right\rceil = \|u\|.
\]
On the other hand, as \( \langle u \rangle \) is principal it follows, according to Lemma 2.3, that if \( I \subseteq \langle u \rangle \) is \( \phi \)-fixed, then \( I \subseteq \langle (\langle u \rangle \cdot m)^\# \rangle \). Nevertheless, it implies that any element \( g \in \langle u \rangle \cdot m \) which forms part of a system of generators for \( \langle u \rangle \cdot m \) is such that
\[
\|g\| \geq \|u\| + 1 > \|u\|.
\]
But this strict lower inequality implies, taking into account the very definition of the hash operation, that \( \langle (\langle u \rangle \cdot m)^\# \rangle = 0 \) and therefore \( I = 0 \), just what we finally wanted to check. \( \square \)

3. Examples

The goal of this section is to present some interesting calculations which were carried out with an implementation of the algorithm presented in this manuscript. Macaulay2 (cf. [2]) has been used extensively both in constructing and exploring examples, as well as implementing the procedure described herein.

Firstly, we include an example where we develop the algorithm step by step for the convenience of the reader.

**Example 3.1.** We consider the ring \( S := \mathbb{F}_2[x, y] \) and set \( u := xy \). We compute \( FP_1(S) \).

(a) Start with \( I = S = I^\# \). As \( I_1(uI) = I \) add \( S \) to the list \( FP_1(S) \).

(b) As \( I \) is principal, go on with \( I = m = I^\# \). Since \( I_1(uI) = I \) add \( m \) to the list \( FP_1(S) \). Moreover, we have to note that
\[
\{ m^2 \subseteq V \subseteq m \mid \dim_{\mathbb{F}_2} m/V = 1 \} = \{ \langle x, y^2 \rangle, \langle y, x^2 \rangle, \langle x^2, xy, x + y \rangle \}.
\]
We have to emphasize that in the calculation of this set is when we are using that we are working with characteristic two.
Thus, we need to compute the following sets of fixed ideals:

$$\text{FP}_1((x^2, xy, x + y)), \text{FP}_1((x, y^2)) \text{ and } \text{FP}_1((y, x^2)).$$

As $\langle x^2, xy, x+y \rangle^\# = \langle xy \rangle$ and $I_1(u(x)) = \langle xy \rangle$ add $\langle xy \rangle$ to the list $\text{FP}_1((x^2, xy, x+y))$. Moreover, as $\langle xy \rangle$ is principal go on with $\langle x^2y, y^2 \rangle$. Nevertheless, since $\langle x^2y, y^2 \rangle^\# = \{0\}$ we deduce that

$$\text{FP}_1((x^2, xy, x+y)) = \{\langle xy \rangle, \langle 0 \rangle\}.$$

On the other hand, since $\langle x, y^2 \rangle^\# = \langle x \rangle$ and $I_1(u(x)) = \langle x \rangle$ add $\langle x \rangle$ to the list $\text{FP}_1((x, y^2))$. In addition, as $\langle x \rangle$ is principal go on with $\langle x^2, xy \rangle$. However, since $\langle x^2, xy \rangle^\# = \langle xy \rangle$ we can use the foregoing calculations and therefore we conclude that $\text{FP}_1((x, y^2)) = \{\langle x \rangle, \langle xy \rangle, \langle 0 \rangle\}$.

A similar computation shows that $\text{FP}_1((y, x^2)) = \{\langle y \rangle, \langle xy \rangle, \langle 0 \rangle\}.$

In this way, it follows that $\text{FP}_1(S) = \{F_2[x, y], \langle x, y \rangle, \langle x \rangle, \langle y \rangle, \langle xy \rangle, \langle 0 \rangle\}.$

Secondly, we include a more involved example which looks in greater detail at [7] Section 9.

**Example 3.2.** Consider the matrix of variables

$$A := \begin{pmatrix} x_1 & x_2 & x_3 & x_4 \\ x_4 & x_2 & x_3 & x_1 \end{pmatrix}$$

and set $S := F_2[x_1, x_2, x_3, x_4].$ Furthermore, for any $1 \leq i < j \leq 4$ we denote by $M_{ij}$ the minor of $A$ of size 2 obtained from columns $i$ and $j$. In addition, set

$$u := x_1^3x_2x_3+x_1^3x_2x_4+x_1^2x_3x_4x_5+x_1x_2x_3x_4x_5+x_1x_2x_4^2x_5+x_2^2x_4x_5+x_3x_4^2x_5+x_3x_4^2x_5^2+x_3^2x_5^2.$$

Our procedure produces the following 84 proper $\phi$-fixed ideals of $S$, where $\phi := u\Phi_1$.

(i) One prime ideal generated by five elements; namely, the ideal $m$ generated by all the variables of $S$.

(ii) Four prime ideals generated by four elements; namely,

$$\langle x_1, x_2, x_3, x_4 \rangle, \langle x_1, x_2, x_4, x_5 \rangle, \langle x_1, x_3, x_4, x_5 \rangle, \langle x_1, x_2, x_3 + x_4, x_5 \rangle.$$

(iii) Five prime ideals generated by three elements; namely,

$$\langle x_1, x_2, x_5 \rangle, \langle x_1, x_3, x_4 \rangle, \langle x_1, x_2, x_4 \rangle, \langle x_1, x_4, x_5 \rangle, \langle x_1 + x_2, x_3 + x_4, x_2^2 + x_4x_5 \rangle.$$

(iv) Two prime ideals generated by two elements; namely, $\langle x_1, x_4 \rangle$ and $\langle x_1 + x_2, x_2^2 + x_4x_5 \rangle$. The reader should notice that

$$\langle x_1 + x_2, x_2^2 + x_4x_5 \rangle = \langle x_1 + x_2, x_1^2 + x_4x_5 \rangle$$

because of we are working on characteristic two.

(v) One prime ideal generated by just one element; namely, the ideal $\langle u \rangle$.

(vi) Twenty-nine ideals which contains in their set of minimal generators some $M_{ij}$ for some $1 \leq i < j \leq 4$.

(vii) The remainder fourty-two ideals define arrangements of linear varieties. Among these 42 ideals, there is one distinguished element; namely, the ideal $\langle x_1, x_2, x_3 + x_4, x_4x_5 \rangle$. In [7] Section 9 it was shown that this ideal is the parameter test ideal of the quotient ring $S/I$, where $I$ is the ideal of $S$ generated by the $2 \times 2$ minors of $A$. 
The reader should notice that, in this case, the set of \( \phi \)-fixed ideals equals the set of \( \phi \)-compatible ideals; indeed, this is due to the fact that, in this case, the map \( \phi = u\Phi \) is a Frobenius splitting. In particular, we recover the thirteen non-zero \( \phi \)-compatible primes obtained by M. Katzman and K. Schwede in [9, Example 7.2].

Thirdly, we include an example where the characteristic of our ground field is greater than two.

**Example 3.3.** Let \( S := \mathbb{F}_5[x, y, z] \), and \( u = x^4 + y^4 + z^4 \). The aim of this example is to compute, using our algorithm, all the \( \phi \)-fixed ideals of \( S \), where \( \phi := u^4\Phi \).

Our method produces the following sixty-five non-zero \( \phi \)-fixed ideals.

(i) The ideal \( \mathfrak{m} := \langle x, y, z \rangle \subseteq S \), its square \( \mathfrak{m}^2 \) and the principal ideal \( \langle u \rangle \).

(ii) Thirty-one ideals of the form \( \mathfrak{m}^2 + \mathfrak{H} \), where \( \mathfrak{H} \) is an ideal of \( S \) generated by a single linear form.

(iii) Thirty-one ideals of the form \( \mathfrak{m}^2 + \mathfrak{G} \), where \( \mathfrak{G} \) is an ideal of \( S \) generated by two linear forms.

It is worth noting that this specific calculation is also interesting because it provides an example where our method provides more information than the procedures worked out in [9]; indeed, if one uses [10] here, then one only gets the ideals \( \langle u \rangle \) and \( \mathfrak{m} \). As we have explained in the Introduction, the reader should remember that, whereas our algorithm produces all the \( \phi \)-fixed ideals, the procedures described in [9] describes algorithmically the radical \( \phi \)-compatible ideals.

Finally, we conclude this paper with the following example, which looks in greater detail at [8, Section 2].

**Example 3.4.** We fix the \( 2 \times 3 \) matrix of indeterminates

\[
\begin{pmatrix}
x_1 & x_2 & x_3 \\
y_1 & y_2 & y_3
\end{pmatrix}
\]

and we let \( S \) to be the polynomial ring \( \mathbb{F}_2[x_1, x_2, x_3, y_1, y_2, y_3] \). Moreover, for any \( 1 \leq i < j \leq 3 \) \( \Delta_{ij} \) will stand for the \( 2 \times 2 \) minor obtained from columns \( i \) and \( j \). In this way, taking into account this notation, we set

\[
u := \Delta_{12}\Delta_{13} = (x_1y_2 - x_2y_1)(x_1y_3 - x_3y_1).
\]

Our procedure produces the following seven proper \( \phi \)-fixed ideals, where \( \phi := u\Phi \);

namely,

\[
\langle x_1, y_1, \Delta_{23} \rangle, \langle x_1, y_1 \rangle, \langle \Delta_{12}, \Delta_{13}, \Delta_{23} \rangle, \langle \Delta_{12}, \Delta_{13} \rangle, \langle \Delta_{12} \rangle, \langle \Delta_{13} \rangle, \langle \Delta_{12}\Delta_{13} \rangle.
\]

In particular, we obtain the following five proper \( \phi \)-fixed prime ideals:

\[
\langle \Delta_{12} \rangle, \langle \Delta_{13} \rangle, \langle x_1, y_1 \rangle, \langle x_1, y_1, \Delta_{23} \rangle, \langle \Delta_{12}, \Delta_{13}, \Delta_{23} \rangle.
\]

Such list of \( \phi \)-fixed prime ideals turns out to be the complete list of proper \( \phi \)-compatible prime ideals, as the reader can check using [10].

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Department of Economics and Business, Universitat Pompeu Fabra, Jaume I Building, Ramon Trias Fargas 25-27, 08005 Barcelona, Spain.
E-mail address: alberto.fernandezb@upf.edu
URL: http://atlas.mat.ub.edu/personals/aboix/

Department of Pure Mathematics, University of Sheffield, Hicks Building, Sheffield S3 7RH, United Kingdom
E-mail address: M.Katzman@sheffield.ac.uk
URL: http://www.katzman.staff.shef.ac.uk/