A Single-Letter Upper Bound to the Mismatch Capacity
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Abstract

We derive a single-letter upper bound to the mismatched-decoding capacity for discrete memoryless channels. The bound is expressed as the mutual information of a transformation of the channel, such that a maximum-likelihood decoding error on the translated channel implies a mismatched-decoding error in the original channel. In particular, a strong converse is shown to hold for this upper bound: if the rate exceeds the upper-bound, the probability of error tends to 1 exponentially when the block-length tends to infinity. We also show that the underlying optimization problem is a convex-concave problem and that an efficient iterative algorithm converges to the optimal solution. In addition, we show that, unlike achievable rates in the literature, the multiletter version of the bound does not improve. A number of examples are discussed throughout the paper.

I. INTRODUCTION AND PRELIMINARIES

We consider reliable communication over a discrete memoryless channel (DMC) \( W \) with a given decoding metric \([1]–[4]\) (see also \([5]\) and references therein for an account of more recent progress). This problem arises when the decoder uses a suboptimal decoding rule due to limited computational resources, simpler implementation, or imperfect channel estimation. Moreover, it is shown in \([1]\) that some important problems in information theory, like the zero-error capacity of a channel can be cast as instances of the mismatch decoding problem. As a result, deriving a single letter characterization of the mismatch decoding capacity would yield a solution to zero-error capacity problem, known to be a difficult problem.

Multiple achievability results have been reported in the literature \([1]–[4], [6]–[8]\). These results were derived by random coding techniques, i.e. analyzing the average probability of error of mismatch decoding over a certain ensemble of codebooks. In some cases, multiuser achievable rates have been shown to improve over standard single-user random coding \([6]–[8]\). As suggested by \([1]\), multiletter versions of achievable rates can yield strict improvements over their single-letter counterparts.

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Unlike the achievable rate case, few converse results have been reported in the literature. The only single-letter converse was reported in [9], where it was claimed that for binary-input DMCs, the mismatch capacity was precisely the LM rate derived in [3], [4]. Reference [10] provided a counterexample to this converse invalidating its claim, showing that a multiletter multiuser rate from [7], [8] was higher than the LM rate. Multiletter converse results were derived in [11]. In particular, for DMCs, [11] shows that for rational decoding metrics, the probability of error cannot decay faster than \( O(n^{-1}) \) for rates above the achievable rate in [3], [4].

In this paper, we propose a single-letter upper bound to the mismatch capacity that is shown to match the mismatch capacity in the cases where it is known, and yield strict improvements over the matched capacity in cases where the mismatch capacity is unknown. The bound is cast as a max-min optimization of the mutual information between the input and a transformation of the channel, such that a maximum-likelihood decoding error on the transformed channel implies a mismatched-decoding error in the original channel. The bound is shown to be convex-concave and an efficient algorithm to compute the bound is provided. The convexity analysis of the bound shows that the multiletter version cannot improve over its single-letter version.

The paper is structured as follows. In Section II we introduce notation and preliminaries. In Section III we introduce our main result and discuss its application to some examples. Sections IV, V, VI, and VII provide the proof of our main result. In particular, in Section IV we construct a graph between different type classes as a key first step of the proof of our upper bound. In Section V we prove a theorem that relates the maximum-likelihood decoding errors on a constructed auxiliary channel \( V \) and mismatched decoding errors on channel \( W \). In Section VI we connect the continuous distributions to the theorems we have derived by the method of types in the previous sections. Section VII gives the final steps of the proof. In Section VIII we show that the optimization problem implied by our bound is a convex-concave optimization problem and we derive the corresponding KKT conditions. Section IX discusses the computation of the bound and proves the convergence of an efficient iterative algorithm. In Section X we use the KKT conditions derived for the single-letter bound and show that the multiletter version of the bound cannot improve over its single-letter counterpart.

II. NOTATION AND PRELIMINARIES

We assume input and output alphabets are \( \mathcal{X} = \{1, 2, \cdots, J \} \) and \( \mathcal{Y} = \{1, 2, \cdots, K \} \), respectively. We denote the channel transition probability by \( W(k|j) \) and define \( W \in \mathbb{R}^{J \times K} \) as the matrix defined by the channel \( W(j, k) = W(k|j) \). A codebook \( C_n \) is defined as a set of \( M \) sequences \( C_n = \{ x(1), x(2), \cdots, x(M) \} \), where \( x(m) = (x_1(m), x_2(m), \cdots, x_n(m)) \in \mathcal{X}^n \), for \( m \in \{1, 2, \cdots, M \} \). A message \( m \in \{1, 2, \cdots, M \} \) is chosen equiprobably and \( x(m) \) is sent over the channel. The channel produces a noisy observation \( y = (y_1, y_2, \cdots, y_n) \in \mathcal{Y}^n \) according to \( W^n(y|x) = \prod_{i=1}^n W(y_i|x_i) \). Upon observing \( y \in \mathcal{Y}^n \) the decoder produces an estimate of the transmitted message \( \hat{m} \in \{1, 2, \cdots, M \} \). The average and maximal error probabilities are respectively defined as

\[
P_e(C_n) = \frac{1}{M} \sum_{i=1}^M \mathbb{P}[\hat{m} \neq m | m = i]
\]

and

\[
P_{e,\text{max}}(C_n) = \max_{i \in \{1, 2, \cdots, M\}} \mathbb{P}[\hat{m} \neq m | m = i].
\]
Rate \( R > 0 \) is said to be achievable if for any \( \epsilon > 0 \) there exists a sequence of length-\( n \) codebooks \( \{C_n\}_{n=1}^{\infty} \) such that \( |C_n| \geq 2^{n(R-\epsilon)} \), and \( \liminf_{n \to \infty} P_e(C_n) = 0 \). The capacity of \( W \), denoted by \( C(W) \) or \( C(W) \), is defined as the largest achievable rate.

The decoder that minimizes the error probability is the maximum-likelihood (ML) decoder, that produces the message estimate \( \hat{m} \) according to

\[
\hat{m} = \arg \max_{i \in \{1, 2, \ldots, M\}} W^n(y|x(i)).
\]  

(3)

In certain situations, it is not possible to use ML decoding and instead, the decoder produces the message estimate \( \hat{m} \) as

\[
\hat{m} = \arg \max_{i \in \{1, 2, \ldots, M\}} d(x(i), y),
\]

(4)

where,

\[
d(x(i), y) = \sum_{\ell=1}^{n} d(x(\ell), y_\ell)
\]

and \( d : \mathcal{X} \times \mathcal{Y} \to \mathbb{R} \) is the decoding metric. We will refer to this decoder as \( d \)-decoder. When \( d(x, y) = \log W(y|x) \), the decoder is ML, otherwise, for a general decoding metric \( d \) the decoder is said to be mismatched.

Lower bounds for the mismatch capacity have been studied extensively using random coding techniques. Specifically, the i.i.d. random coding ensemble is known to achieve the generalized mutual information (GMI) which can be written as

\[
R_{d}^{\text{GMI}}(W) = \max_{P_X} \min_{P_{X|Y|X}, P_Y|X} \mathbb{E}_{P_X} \left[ I(P_X, V_{Y|X}) \right] \]

An improved lower bound, known as the LM rate, is derived by employing constant composition random coding.

\[
R_{d}^{\text{LM}}(W) = \max_{P_X} \min_{P_{X|Y|X}, P_Y|X} \mathbb{E}_{P_X} \left[ I(P_X, V_{Y|X}) \right] \]

The above rate has an intuitive explanation. The maximization is over all input distributions, and the minimizations is over all auxiliary channels \( V_{Y|X} \) with two properties. First, equal output marginal \( P_X \times V_{Y|X} = P_X \times W \), such that for all \( k \in \mathcal{Y} \)

\[
\sum_{j \in \mathcal{X}} P_X(j)V_{Y|X}(k|j) = \sum_{j \in \mathcal{X}} P_X(j)W(k|j).
\]  

(8)

This implies that the distribution of the received sequence needs to be the same for both channel \( W \) and auxiliary channel \( V_{Y|X} \) whenever the input codeword is chosen from composition \( P_X \). The second condition \( \mathbb{E}_{P_X \times V_{Y|X}}[d(X, Y)] \geq \mathbb{E}_{P_X \times W}[d(X, Y)] \) can be rewritten as,

\[
\sum_{j, k} P_X(j)V_{Y|X}(k|j)d(j, k) \geq \sum_{j, k} P_X(j)V_{Y|X}(k|j)d(j, k) \]

(9)

and implies that, the received sequence \( Y \) has a higher metric under channel \( V_{Y|X} \) than under channel \( W \), and therefore, the \( d \)-decoder makes an error. It is implied in \[6\] and \[7\] that \( R_{d}^{\text{GMI}}(W) \leq R_{d}^{\text{LM}}(W) \). The GMI and
LM rates are ensemble tight, i.e. the ensemble average error probability tends to one exponentially for rates exceeding the GMI and LM rates, respectively. Both of the bounds above are known not to be tight in general.

It is known that the GMI and LM rates can be improved by considering their multiletter counterparts [1]. The method of types [13] Ch. 2 will be used extensively in this paper. We recall some of the basic definitions and introduce some notation. The type of a sequence \( x = (x_1, x_2, \cdots, x_n) \in \mathcal{X}^n \) is a column vector representing its empirical distribution, i.e., \( \hat{p}_x(j) = \frac{1}{n} \sum_{i=1}^{n} 1 \{x_i = j\} \). The set of all types of \( \mathcal{X}^n \) is denoted by \( \mathcal{P}_n(\mathcal{X}) \).

For \( p_X \in \mathcal{P}_n(\mathcal{X}) \), the type class \( T^n(p_X) \) is set of all sequences in \( \mathcal{X}^n \) with type \( p_X \). \( T^n(p_X) = \{ x \in \mathcal{X}^n | \hat{p}_x = p_X \} \). Throughout the paper, for two \( J \times K \) matrices \( M_1, M_2 \) representing conditional types we define

\[
|M_1 - M_2|_\infty = \max_{1 \leq j \leq J} \sum_{k=1}^{K} |M_1(j,k) - M_2(j,k)|. \tag{10}
\]

The joint type of sequences \( x = (x_1, x_2, \cdots, x_n) \in \mathcal{X}^n \) and \( y = (y_1, y_2, \cdots, y_n) \in \mathcal{Y}^n \) is defined as a matrix representing their empirical distribution \( \hat{p}_{xy}(j,k) = \frac{1}{n} \sum_{i=1}^{n} 1 \{x_i = j, y_i = k\} \). The conditional type of \( y \) given \( x \) is the matrix

\[
\hat{p}_{y|x}(j,k) = \begin{cases} 
\hat{p}_{xy}(j,k) / \hat{p}_x(j) & \hat{p}_x(j) > 0 \\
0 & \text{otherwise.}
\end{cases} \tag{11}
\]

The set of all joint conditional types on \( \mathcal{Y}^n \) given \( \mathcal{X}^n \) is denoted by \( \mathcal{P}_n(\mathcal{Y}|\mathcal{X}) \). For \( p_{Y|X} \in \mathcal{P}_n(\mathcal{Y}|\mathcal{X}) \) and sequence \( x \in T^n(p_X) \), the conditional type class \( T^n_{y|x}(p_{Y|X}) \) is defined as \( T^n_{y|x}(p_{Y|X}) = \{ y \in \mathcal{Y}^n | \hat{p}_{y|x} = p_{Y|X} \} \).

Similarly, we can define the joint type of \( x, y, \hat{y} \) as the empirical distribution of the triplet. For \( j \in \mathcal{X} \) and \( k_1, k_2 \in \mathcal{Y} \),

\[
\hat{p}_{x,y,\hat{y}}(j,k_1,k_2) = \frac{1}{n} \sum_{i=1}^{n} 1 \{x_i = j, y_i = k_1, \hat{y}_i = k_2\}. \tag{12}
\]

We define the joint conditional type of \( y, \hat{y} \) given \( x \in T^n(p_X) \) as

\[
\hat{p}_{y,\hat{y}|x}(j,k_1,k_2) = \begin{cases} 
\hat{p}_{x,y,\hat{y}}(j,k_1,k_2) / \hat{p}_x(j) & \hat{p}_x(j) > 0 \\
0 & \text{otherwise.}
\end{cases} \tag{13}
\]

The set of all joint conditional types is denoted by \( \mathcal{P}_n(\mathcal{Y}|\mathcal{Y},\mathcal{X}) \). Additionally, for \( p_{Y|Y|X} \in \mathcal{P}_n(\mathcal{Y}|\mathcal{Y}|\mathcal{X}) \) we define:

\[
T^n_{y|x}(p_{Y|Y|X}) = \{ \hat{y} \in \mathcal{Y}^n | \hat{p}_{y|x} = p_{Y|Y|X} \} \tag{14}
\]

The mutual information is defined as

\[
I(p_X, p_{Y|X}) \triangleq \mathbb{E} \left[ \log \frac{p_{Y|X}(y|x)}{\sum_{y'} p_{X}(x)p_{Y|X}(y'|x)} \right].
\]

**Definition 1:** Let \( P_{Y|Y|X} \) be a joint conditional distribution and define the set

\[
S_d(k_1,k_2) \triangleq \{ i \in \mathcal{X} | i = \arg \max_{i' \in \mathcal{X}} d(i', k_2) - d(i', k_1) \}. \tag{15}
\]

We say that \( P_{Y|Y|X} \) is a maximal joint conditional distribution if for all \( (j, k_1, k_2) \in \mathcal{X} \times \mathcal{Y} \times \mathcal{Y} \),

\[
P_{Y|Y|X}(j, k_1, k_2) = 0 \quad \text{if} \quad j \notin S_d(k_1,k_2). \tag{16}
\]

Moreover, if \( p_{Y|Y|X} \in \mathcal{P}_n(\mathcal{Y}|\mathcal{Y}|\mathcal{X}) \) satisfies the same condition, we call it a maximal joint conditional type.

For a given decoding metric \( d \), we define the set of maximal joint distributions to be \( \mathcal{M}_{\text{max}}(d) \).
The above definition will become helpful when relating decoding errors in channel $P_{Y|X} = W$ under $d$-decoding to errors in channel $P_{Y|X}$ under ML decoding.

**Definition 2:** Let $C_n = \{x(1), \cdots, x(M)\}$ and $m$ be the transmitted message. We say that the decoder makes a **type conflict error** for a given $y \in Y^n$ if there is at least one codeword $x(i) \neq x(m)$ such that that $\hat{P}_{y|x(i)} = \hat{P}_{y|x(m)}$.

If there is a type conflict error, every decoder that makes a decision based on the joint type between the channel output and the candidate codewords ($\alpha$-decoder) makes an error, including ML and $d$-decoding; the converse is not true. With the same method developed in the paper, it can be shown that the type conflict error probability over channel $W$ goes to 1 exponentially for $R > C(W)$; even with a genie-aided ML decoder knowing the exact conditional type $\hat{P}_{y|x(m)}$, the error probability would still tend to 1 exponentially above capacity.

### III. MAIN RESULT

In this section, we introduce the main result and discuss some of its properties.

**Theorem 1:** Let $W, d$ be channel and decoding metric, respectively. We define $\bar{R}_d(W)$ as follows

$$
\bar{R}_d(W) = \max_{P_X} \min_{P_{Y|X} \in M_{\text{max}}(d)} I(P_X, P_{Y|X}).
$$

If $R > \bar{R}_d(W)$, $\exists n_0 \in \mathbb{N}$ and $E_{\text{sc}}^d(R) > 0$ such that for $n > n_0$, the error probability of codebook $C_n$ of length $n$ and $M \geq 2^{nR}$ codewords satisfies $P_{e_{\text{max}}}(C_n) \geq 1 - 2^{-nE_{\text{sc}}^d(R)}$.

**Proof:** The proof is developed over the next sections of the paper. The main idea behind proof of Theorem 1 is that of lower-bounding the error probability of a codebook $C_n$ with $d$-decoding over channel $W$ by that of the same codebook over a different channel $V$ with ML decoding, with $V = P_{Y|X}$ as per the theorem statement. We will be able to construct a graph $G$ in the output space such that if ML decoding over $V$ makes a type conflict error for some $y \in Y^n$, then, the $d$-decoder makes an error for some $\hat{y} \in Y^n$ connected to $y$ in $G$.

Theorem 1 implies that $C_d(W) \leq \bar{R}_d(W)$. It is implied in Theorem 1 that for any $P_{Y|X} \in M_{\text{max}}(d)$ such that $P_{Y|X} = W$,

$$
\bar{R}_d(W) \leq C(P_{Y|X}).
$$

This result is derived by using the max-min inequality:

$$
\bar{R}_d(W) = \max_{P_X} \min_{P_{Y|X} \in M_{\text{max}}(d)} I(P_X, P_{Y|X})
\leq \min_{P_{Y|X} \in M_{\text{max}}(d)} \max_{P_X} I(P_X, P_{Y|X})
= \min_{P_{Y|X} \in M_{\text{max}}(d)} C(P_{Y|X}).
$$
construct bounds without necessarily performing the optimization. As an instance of the above result, setting $Y$ such that $P_{Y|X} = W$ and $Y = Y$ makes $P_{Y|X}$ a maximal joint conditional distribution (Def. [1]). Therefore, $C_d(W) \leq C(P_{Y|X}) = C(W)$. Moreover, the bound above can be used for a specific input distribution. In the proof it is evident that the bound remains valid for any fixed input distribution, not only the maximizing one.

Remark 1: The optimization (17) in Theorem [1] is a convex-concave optimization problem. See Section VIII for further details.

Remark 2: It was shown [1] that the achievability bounds for DMCs could be improved by considering an equivalent metric $d(x, y) = sd(x, y) + a(x) + b(y)$. Here we show that our bound in [1] does not change by replacing metric $d(x, y)$ by $d(x, y) = sd(x, y) + a(x) + b(y)$. According to the definition of $\mathcal{S}_d(k_1, k_2)$, we have

$$\arg\max_{j \in X} d(j, k_2) - d(j, k_1) = \arg\max_{j \in X} (sd(j, k_2) + a(j) + b(k_2)) - (sd(j, k_1) + a(j) + b(k_1))$$

which is precisely the condition in the definition of $S_d(k_1, k_2)$.

The above property from [1] suggests that for binary-input channels, the mismatch capacity $C_d(W)$ is only a function of the metric difference $d(1, y) - d(2, y)$. In the remainder of this section, we show a sufficient condition for $C_d(W) < C(W)$ for binary-input channels based on the above observation.

Definition 3: We say that two sequences $\{\alpha_i\}_{i=1}^K$ and $\{\beta_i\}_{i=1}^K$ have the same order if for all $1 \leq i_1, i_2 \leq K$

$$\alpha_{i_1} \geq \alpha_{i_2} \Rightarrow \beta_{i_1} \geq \beta_{i_2}. \quad (25)$$

We have the following result for $|X| = 2$.

Theorem 2: Assume that $W(k|j) > 0$, for all $j = 1, 2$, $k = 1, \ldots, K$. If the sequences $\{\log W(k|1) - \log W(k|2)\}_{k=1}^K$ and $\{d(1, k) - d(2, k)\}_{k=1}^K$ do not have the same order, then $C_d(W) < C(W)$.

Proof: See Appendix for the proof.

A. Examples

In the following, we discuss the applicability of our upper bound to two relevant cases. First, we show that our bound recovers known results on binary-input binary-output channels. Next, we show that our bound makes a non-trivial improvement over the channel-metric combination used in [10] to state the counterexample to Balakirsky’s result [9].

Example 1 (Binary-input binary-output channels): Suppose that the channel and decoding metric matrices of binary-input binary-output channels are given by

$$W = \begin{bmatrix} a & b \\ c & d \end{bmatrix} \quad \text{and} \quad D = \begin{bmatrix} \hat{a} & \hat{b} \\ \hat{c} & \hat{d} \end{bmatrix}. \quad (26)$$

Without loss of generality we assume $a + d \geq b + c$. We show the following known result [1]: if $\hat{a} + \hat{d} < \hat{b} + \hat{c}$ then $\bar{R}_d(W) = 0$. On the other hand, if $\hat{a} + \hat{d} \geq \hat{b} + \hat{c}$, then $\bar{R}_d(W) = C(W)$.

Case 1: $\hat{a} + \hat{d} < \hat{b} + \hat{c}$
We chose the joint conditional distribution in Table I

It can be checked that indeed it is a valid joint conditional distribution for $0 \leq r_1 \leq a$ and $0 \leq r_2 \leq d$, and that $\sum_{k_2} P_{Y|X}(j, k_1, k_2) = P_{Y|X}(j, k_1) = W(k_1(j))$. In order to check its maximality, we first notice that for $k_1 = k_2$ we always have that $d(i, k_2) - d(i, k_1) = 0$ for all $i \in X$, implying that $S_d(k_1, k_2) = \{1, 2\}$. Thus, since every $j \in X$ is such that $j \in S_d(k_1, k_2)$, the corresponding four entries can be nonzero. As for entry $(1, 1, 2)$ (resp. $(2, 2, 1)$), using the assumption $\hat{a} + \hat{d} < \hat{b} + \hat{c}$ we have that $S_d(k_1, k_2) = \{1\}$ (resp. $S_d(k_1, k_2) = \{2\}$), and thus they both can be nonzero. Since by assumption $\hat{a} + \hat{d} < \hat{b} + \hat{c}$, it can be checked that for entry $(2, 1, 2)$, $S_d(k_1, k_2) = \{1\}$, and thus we must have $P_{Y|X}(j, k_1, k_2) = 0$. Similarly for entry $(1, 2, 1)$, $S_d(k_1, k_2) = \{2\}$. Marginalizing the above over $Y$ gives

$$P_{Y|X} = \begin{bmatrix} a - r_1 & b + r_1 \\ c + r_2 & d - r_2 \end{bmatrix}. \quad (27)$$

Without loss of generality assume that $a$ is the largest element of $W$. By setting $r_1 = r_2 = \frac{a - c}{2} = \frac{d + b}{2}$ we obtain

$$P_{Y|X} = \begin{bmatrix} \frac{a + c}{2} & \frac{b + d}{2} \\ \frac{a + c}{2} & \frac{b + d}{2} \end{bmatrix}. \quad (28)$$

Since $C(P_{Y|X}) = 0$, we have that $C_d(W) \leq 0$.

**Case 2:** $\hat{a} + \hat{c} \geq \hat{b} + \hat{c}$

In [9] it is shown that the LM achievable rate is equal to $C(W)$. Therefore, our upper-bound also matches the achievable rate.

**Example 2:** We consider the channel and metric studied in [10] to show a counterexample to [9]

$$W = \begin{bmatrix} 0.97 & 0.03 & 0 \\ 0.1 & 0.1 & 0.8 \end{bmatrix} \quad \text{and} \quad D = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 0.5 & 1.36 \end{bmatrix}. \quad (29)$$

In this case, the LM rate is $R_d^{LM}(W) = 0.1975$ while the rate achieved by a multiletter extension of order $\ell = 2$ of superposition coding gives $R_d^{SC,(2)}(W) = 0.1991$ [10].

We choose the maximal $P_{Y|X}$ in Table I such that $P_{Y|X} = W$, which happens to be the optimal one (see Section IX for details). By marginalizing over $Y$ we find that

$$P_{Y|X} = \begin{bmatrix} 0.5 & 0.5 & 0 \\ 0.1 & 0.1 & 0.8 \end{bmatrix}. \quad (30)$$

We obtain that $\bar{R}_d(W) = 0.6182 \text{ bits/use}$, while the capacity is $C(W) = 0.7133 \text{ bits/use}$.
In the above example, if we change \(d(2, 2)\) from 0.5 to 1, the same \(P_{Y|X}^j\) in Table II remains maximal (and optimal) and gives \(\hat{R}_d(W) = 0.6182\) bits/use, matching the LM rate \([4]\).

**Example 3:** In this example we apply our bound to the erasures-only or zero-undetected error capacity problem. In this setting, the decoder chooses a codeword \(x\) in the codebook if it is the only codeword with \(W(y|x) > 0\). Otherwise the decoder declares an erasure. The erasures-only capacity \(C_{eo}(W)\) is defined as the maximum achievable rate where the probability of erasure could tend to zero by increasing the block-length. It can be shown \([11]\) that the erasures-only capacity problem can be reduced to a mismatched decoding problem with decoding metric

\[
d(x, y) = \begin{cases} 
0 & W(y|x) > 0 \\
-1 & W(y|x) = 0.
\end{cases}
\]  

(31)

Now for any two fixed outputs \(k_1, k_2 \in \mathcal{Y}\), there are a few different possibilities. First, if there is an input \(j \in \mathcal{X}\) such that, \(W(k_1|j) = 0\) and \(W(k_2|j) > 0\) then clearly \(j \in S_d(k_1, k_2)\). Thus, \(P_{Y|X}^j(k_1, k_2|j)\) could potentially be non-zero. Yet, since \(P_{Y|X}^j(k_1, k_2|j) \leq W(k_1|j)\), then, \(P_{Y|X}^j(k_1, k_2|j) = 0\) if \(W(k_1|j) = 0\). Instead, if there is no input \(j \in \mathcal{X}\) such that \(W(k_1|j) = 0\) and \(W(k_2|j) > 0\), then

\[
\{j \in \mathcal{X}|W(k_1|j) > 0\} \subseteq \{j \in \mathcal{X}|W(k_2|j) > 0\}. 
\]  

(32)

If \(\{j \in \mathcal{X}|W(k_1|j) > 0\} = \{j \in \mathcal{X}|W(k_2|j) > 0\}\), then, outputs \(k_1\) and \(k_2\) can be merged without affecting \(C_{eo}(W)\) \([14]\). Otherwise, outputs \(k_1\) and \(k_2\) cannot be merged.

Consider now the following ternary-input quaternary-output channel that cannot be simplified by merging,

\[
W = \begin{bmatrix}
0.25 & 0 & 0.05 & 0.7 \\
0.3 & 0.55 & 0 & 0.15 \\
0.05 & 0.5 & 0.45 & 0
\end{bmatrix}.
\]  

(33)

The Shannon capacity of \(W\) is \(C(W) = 0.7854\) bits/use and our upper bound gives \(\hat{R}_d(W) = 0.6232\) bits/use.

The LM rate computed by an exhaustive search over the input distributions is \(R_{LM}^d(W) = 0.4292\) bits/use.

As observed from the numbers, in all cases our bound non-trivially improves on the Shannon capacity.

**IV. Graph Construction**

In this section, we outline how to construct a graph between two different conditional types obtained from a joint conditional type.
**Definition 4:** Let $\mathcal{G} = \{\mathcal{V}_1, \mathcal{V}_2, \mathcal{E}\}$ be a regular bipartite graph with vertex sets $\mathcal{V}_1$ and $\mathcal{V}_2$, edge set $\mathcal{E}$ and degrees $r_1$ on vertex set $\mathcal{V}_1$ and $r_2$ on vertex set $\mathcal{V}_2$. For $B \subset \mathcal{V}_2$ we define the set of vertices in $\mathcal{V}_1$ connected to $B$ as

$$\Psi_{21}(B) = \{v \in \mathcal{V}_1 \mid \exists b \in B; (b, v) \in \mathcal{E}\}. \tag{34}$$

Analogously for $B \subset \mathcal{V}_1$, the set $\Psi_{12}(B)$ is defined similar to (34).

**Lemma 1:** Suppose $\mathcal{G} = \{\mathcal{V}_1, \mathcal{V}_2, \mathcal{E}\}$ is a regular bipartite graph with degrees $r_1 > 0, r_2 > 0$. Then, for any $B \subset \mathcal{V}_2$ we have that

$$\frac{|\Psi_{21}(B)|}{|\mathcal{V}_1|} \geq \frac{|B|}{|\mathcal{V}_2|}. \tag{35}$$

**Proof:** Let $B \subset \mathcal{V}_2$ and consider $\Psi_{21}(B)$. There are exactly $r_2|B|$ edges between $B$ and $\Psi_{21}(B)$. Since each vertex in $\Psi_{21}(B)$ is connected to at most $r_1$ vertices of $B$ we have

$$r_2|B| \leq r_1|\Psi_{21}(B)| \tag{36}$$

which implies that

$$\frac{r_2}{r_1}|B| \leq |\Psi_{21}(B)|. \tag{37}$$

Since there are exactly $r_1|\mathcal{V}_1| = r_2|\mathcal{V}_2|$ edges in the graph, the result follows by substituting $\frac{r_2}{r_1} = \frac{|\mathcal{V}_1|}{|\mathcal{V}_2|}$ in (37).

Our aim is to construct a graph between different two conditional type classes, in order to be able to relate type conflict errors of codebook $\mathcal{C}_n$ over channel $V$ and errors of $\mathcal{C}_n$ over channel $W$ under $d$-decoding. Suppose $p_{Y,Y|X} \in \mathcal{P}_n(Y,Y|X)$ is an arbitrary joint conditional type. We construct a graph between $T^n_w(p_{Y|X})$ and $T^n_w(p_{Y|X})$, the corresponding conditional type classes.

**Definition 5:** The graph

$$\mathcal{G}_x(p_{Y|X}) = \{T^n_w(p_{Y|X}), T^n_w(p_{Y|X}), \mathcal{E}\} \tag{38}$$

has the following edge set:

$$\mathcal{E} = \{(y, \hat{y}) | \hat{y} \neq y \mid \hat{y} \notin \text{output of an arbitrary channel } V \} \tag{39}$$

**Lemma 2:** The graph $\mathcal{G}_x(p_{Y|X})$ is regular.

**Proof:** For a given $x \in T^n_w(p_{X})$, $|T^n_w(p_{Y|X})|$ is independent of the chosen $x \in T^n_w(p_{X})$, but dependent on $p_X$. Similarly, for a given $y \in T^n_w(p_{Y|X})$, $|T^n_w(p_{Y|X})|$ is independent of the chosen $x, y$, but dependent on the joint type $p_{XY}$. Therefore, the total number of edges that are connected to any given $y \in T^n_w(p_{Y|X})$ is equal to $|T^n_w(p_{Y|X})|$ (see (44)). This proves the left-regularity, i.e., for vertex set $T^n_w(p_{Y|X})$. The same argument holds for $\hat{y} \in T^n_w(p_{Y|X})$ and therefore the graph is regular.

As we show next, the combination of Lemmas 1 and 2 will prove to be helpful. Assume for a codeword $x$ we find a set $B \subset T^n_w(p_{Y|X})$ that yields to type conflict error. Then, the probability of an element $\hat{y} \in \mathcal{B}$ being the output of an arbitrary channel $V$ given that the conditional type is $p_{Y|X}$, is given by

$$\mathbb{P}[\hat{y} \in \mathcal{B} \mid \hat{y} \in T^n_w(p_{Y|X}), x \text{ is sent}] = \frac{|B|}{|T^n_w(p_{Y|X})|}. \tag{40}$$

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where the probability is computed \( w.r.t. \) an auxiliary memoryless channel \( V \), i.e., \( \mathbb{P} [ \hat{y} | x ] \) is sent = \( \prod_{i=1}^{n} V(\hat{y}_i | x_i) \) and equality holds because all elements of \( T_{\alpha}^n(\mathcal{P}_{Y|X}) \) are equally likely to appear at the output when \( x \) is sent.

Now if the graph \( G_x(\mathcal{P}_{Y|X}) \) is connecting a type conflict error to a \( d \)-decoder error, by Lemma 1 we show that the set \( \Psi_2(\mathcal{B}) \subset T_{\alpha}^n(\mathcal{P}_{Y|X}) \) satisfies

\[
\frac{|\Psi_2(\mathcal{B})|}{|T_{\alpha}^n(\mathcal{P}_{Y|X})|} \geq \frac{|\mathcal{B}|}{|T_{\alpha}^n(\mathcal{P}_{Y|X})|}. \tag{41}
\]

Now by the same argument as in (40) we have

\[
\mathbb{P} [ y \in \Psi_2(\mathcal{B}) \mid y \in T_{\alpha}^n(\mathcal{P}_{Y|X}), x \text{ is sent} ] = \frac{|\Psi_2(\mathcal{B})|}{|T_{\alpha}^n(\mathcal{P}_{Y|X})|}. \tag{42}
\]

Combining (42) and (41) we get

\[
\mathbb{P} [ y \in \Psi_2(\mathcal{B}) \mid y \in T_{\alpha}^n(\mathcal{P}_{Y|X}), x \text{ is sent} ] \geq \mathbb{P} [ \hat{y} \in \mathcal{B} \mid \hat{y} \in T_{\alpha}^n(\mathcal{P}_{Y|X}), x \text{ is sent} ].
\]

As a result, we get a lower bound on the probability of error of the \( d \)-decoder in channel \( W \) as a function of type conflict errors in channel \( V \). In the next section we prove that a graph constructed based on a maximal joint conditional type has the property of connecting type conflict errors to \( d \)-decoder errors.

\[ \text{V. CONNECTING ERRORS} \]

We next introduce a property of maximal joint conditional types and use it to relate type conflict and \( d \)-decoding errors.

\[ \text{Lemma 3:} \] Let \( \mathcal{P}_X \in \mathcal{P}_n(\mathcal{X}), x, \hat{x} \in T^n(\mathcal{P}_X) \), and \( \mathcal{P}_{Y|X} \) be a maximal joint conditional type. If \( \hat{y} \in T_{\alpha}^n(\mathcal{P}_{Y|X}) \cap T_{\alpha}^n(\mathcal{P}_{Y|X}) \) is connected to \( y \in T_{\alpha}^n(\mathcal{P}_{Y|X}) \) in \( G_x(\mathcal{P}_{Y|X}) \) then,

\[
d(\hat{x}, y) \leq d(\hat{x}, y). \tag{43}
\]

\[ \text{Proof:} \] From the definition of type, for any \( \hat{x} \in \mathcal{X}^n,
\]

\[
\hat{p}_{uy}(k_1, k_2) = \sum_j \hat{p}_{xuy}(j, k_1, k_2). \tag{44}
\]

Now, if we use the above equation once by setting \( \hat{x} = x \) and once by setting \( \hat{x} = \hat{x} \) we get

\[
\sum_j \hat{p}_{xuy}(j, k_1, k_2) = \sum_j \hat{p}_{xuy}(j, k_1, k_2). \tag{45}
\]

We continue by bounding \( d(\hat{x}, \hat{y}) \) as

\[
d(\hat{x}, \hat{y}) = n \sum_{j,k} \hat{p}_{xy}(j,k)(d(j,k) - d(j,k)) \tag{46}
\]

\[
\leq n \sum_{j,k} \left( \sum_j \hat{p}_{xuy}(j,k) \right) \max_j (d(j,k) - d(j,k)) \tag{47}
\]

\[
= n \sum_{j,k} \left( \sum_j \hat{p}_{xuy}(j,k) \right) \max_j (d(j,k) - d(j,k)) \tag{48}
\]

\[
= n \sum_{j,k} \sum_j \hat{p}_{xuy}(j,k) \tag{49}
\]

\[
= d(x, \hat{y}) - d(x, y) \tag{50}
\]

where (46) follows from the definition of metric and type, since for a joint type \( \hat{p}_{xy} \) we have that \( d(x, y) = n \sum_{j,k} \hat{p}_{xy}(j,k) d(j,k) \), (47) follows from upper-bounding \( (d(j,k) - d(j,k)) \) by \( \max_j (d(j,k) - d(j,k)) \).
follows from (45). (49) follows from the maximality of \( \hat{p}_{Y|X} \) (see Definition (1)) and graph construction \( G_\epsilon(p_{Y|X}) \) (see Definition (5)) and (50) follows again from the metric definition.

Now, using the fact that \( \hat{y} \in \mathcal{T}_\epsilon^*(p_{Y|X}) \cap T_\epsilon^*(p_{Y|X}) \) we get a type conflict error, i.e., \( \hat{p}_{\hat{y}|x} = \hat{p}_{\hat{y}|\hat{x}} \). Thus, \( d(x, \hat{y}) = d(\hat{x}, \hat{y}) \). Finally, combining with (50) we get the desired result \( d(x, y) \leq d(\hat{x}, \hat{y}) \), i.e., a \( d \)-decoding error.

The condition of the Lemma says that if \( \hat{y} \in \mathcal{T}_\epsilon^*(p_{Y|X}) \cap T_\epsilon^*(p_{Y|X}) \) and if \( x, \hat{x} \in C_n \), by observing \( \hat{y} \) when \( x \) is sent, there will be a type conflict error. Moreover, if such a \( \hat{y} \) is connected to \( y \) in \( G_\epsilon(p_{Y|X}) \), then, based on (43), by observing \( y \) when \( x \) is sent, the \( d \)-decoder makes an error.

Definition 6: Let \( W \) be a channel and \( p_X \in \mathcal{P}_n(X) \) an input type. We define the channel type neighborhood as the set of conditional types that are close to \( W \).

\[
N_{e,p_X}(W) = \{ p_{Y|X} \in \mathcal{P}_n(Y|X) \mid \forall j, k \text{ if } p_X(j) > 0, |W(k|j) - p_{Y|X}(k|j)| \leq \epsilon \}. \tag{51}
\]

The previous result showed that if \( p_{Y|X} \) is a maximal joint conditional type, then type conflict errors in \( \mathcal{T}_\epsilon^*(p_{Y|X}) \) can be related to \( d \)-decoding errors in \( T_\epsilon^*(p_{Y|X}) \). The lemma below shows that for every empirical conditional type \( \hat{W} \) close to the channel \( W \) there exists a maximal joint conditional type, that can be used to relate type conflict errors of a type close to \( V \) to \( d \)-decoder errors over \( W \).

Lemma 4: Let \( p_X \in \mathcal{P}_n(X) \) be an input type and \( p_{\text{min}} \triangleq \min_{j,p_X(j)>0} p_X(j) \). Assume \( p_{Y|X} \) is a maximal joint conditional distribution such that \( p_{Y|X} = W \) and \( p_{Y|X} = V \). Moreover, let \( \epsilon \geq \frac{2|Y|}{n p_{\text{min}}} \). Then for each \( \hat{W} \in N_{e,p_X}(W) \) we can find a maximal joint conditional type \( p_{Y|X} \) such that \( p_{Y|X} = \hat{W} \) and \( p_{Y|X} \in N_{2K\epsilon,p_X}(V) \).

Proof: If \( \epsilon > 1 \) then there is nothing to prove. So let \( \epsilon < 1 \). For \( j \in X \) and \( k_1, k_2 \in Y \), choose \( p_{Y|X}(j, k_1, k_2) \) to be either

\[
p_{Y|X}(j, k_1, k_2) = \frac{np_X(j)p_{Y|X}(j, k_1, k_2)}{np_X(j)} \tag{52}
\]

or

\[
p_{Y|X}(j, k_1, k_2) = \frac{np_X(j)p_{Y|X}(j, k_1, k_2)}{np_X(j)} \tag{53}
\]

such that for every \( j \in X \) we have

\[
\sum_{k_1, k_2} p_{Y|X}(j, k_1, k_2) = 1. \tag{54}
\]

Such a choice is possible since

\[
\sum_{k_1, k_2} p_{Y|X}(j, k_1, k_2) = 1. \tag{55}
\]

Moreover, when \( p_X(j) = 0 \) define \( p_{Y|X}(j, k_1, k_2) \) as in (13). The above choice implies that for every \( j \in X \) such that \( p_X(j) > 0 \) and any \( k_1, k_2 \in Y \),

\[
|p_{Y|X}(j, k_1, k_2) - p_{Y|X}(j, k_1, k_2)| \leq \frac{1}{np_X(j)} \tag{56}
\]

\[
\leq \frac{1}{np_{\text{min}}}. \tag{57}
\]
Moreover, based on (52) and (53) \( p_{Y\mid X} \) is maximal, since \( p_{Y\mid X}(j, k_1, k_2) \) is non-zero either when \( k_1 = k_2 \) or for the same entries that \( P_{Y\mid X} \) is non-zero. As a result of (57) for every \( j \in \mathcal{X} \) such that \( p_X(j) > 0 \) we have that
\[
\left| \sum_{k_2} (p_{Y\mid X}(j, k_1, k_2) - P_{Y\mid X}(j, k_1, k_2)) \right| \leq \sum_{k_2} \left| p_{Y\mid X}(j, k_1, k_2) - P_{Y\mid X}(j, k_1, k_2) \right| \quad (58)
\]
\[
\leq \frac{|\mathcal{Y}|}{np_{\min}} \quad (59)
\]
and thus,
\[
p_{Y\mid X} \in \mathcal{N}_{\frac{|\mathcal{Y}|}{np_{\min}}} p_X(W)
\]
\[
p_{Y\mid X} \in \mathcal{N}_{\frac{|\mathcal{Y}|}{np_{\min}}} p_X(V)
\]

For any \( \mathbf{W} \in \mathcal{N}_{p_X}(W) \), for every \( j \in \mathcal{X} \) such that \( p_X(j) > 0 \) and \( k \in \mathcal{Y} \) by definition we know that \( |\mathbf{W}(j, k) - p_{Y\mid X}(j, k)| \leq \epsilon \), since
\[
|\mathbf{W}(j, k) - p_{Y\mid X}(j, k)| \leq |\mathbf{W}(j, k) - W(j, k)| + |W(j, k) - p_{Y\mid X}(j, k)| \quad (62)
\]
\[
\leq \frac{\epsilon}{2} + \frac{\epsilon}{2} \quad (63)
\]
\[
= \epsilon. \quad (64)
\]

where (63) follows from (60) and (61). Now, construct \( p'_{Y\mid X} \) from \( p_{Y\mid X} \) in the following way. For any \( j, k \in \mathcal{X} \times \mathcal{Y} \) such that \( p_X(j) > 0 \), if \( \mathbf{W}(j, k) - W(j, k) > 0 \) add non-negative real numbers less than or equal \( \epsilon \) to \( p_{Y\mid X}(j, k_1, k_2) \), \( k_2 = 1, 2, \ldots, K \) to obtain \( p'_{Y\mid X} \) with the following property
\[
\sum_{k_2} p'_{Y\mid X}(j, k_1, k_2) = \mathbf{W}(k_1|j). \quad (65)
\]
We can do this because \( |\mathbf{W}(j, k) - p_{Y\mid X}(j, k)| \leq \epsilon \). We can do the same if \( \mathbf{W}(j, k) - W(j, k) \leq 0 \) with non-positive real numbers not less than \( -\epsilon \) such that
\[
\sum_{k_2} p'_{Y\mid X}(j, k_1, k_2) = \mathbf{W}(k_1|j) \quad (66)
\]
Then marginalizing over \( \hat{Y} \) we get \( p'_{Y\mid X} \)
\[
|p'_{Y\mid X}(j, k_2) - p_{Y\mid X}(j, k_2)| = \left| \sum_{k_1} p'_{Y\mid X}(j, k_1, k_2) - p_{Y\mid X}(j, k_1, k_2) \right| \quad (67)
\]
\[
\leq \sum_{k_1} \left| p'_{Y\mid X}(j, k_1, k_2) - p_{Y\mid X}(j, k_1, k_2) \right| \quad (68)
\]
\[
\leq \sum_{k_1} \epsilon \quad (69)
\]
\[
= K \epsilon. \quad (70)
\]

Therefore by the triangle inequality and (61) we get \( |p'_{Y\mid X}(j, k_2) - p_{Y\mid X}(j, k_2)| \leq 2K \epsilon \).

In the next theorem, we show that if \( P_{Y\mid X} \) is a maximal joint conditional distribution and \( M \) is large enough, then we will find many type conflict errors over conditional types close to \( V = P_{Y\mid X} \). These are then linked to \( d \)-decoding errors over channel \( W = P_{Y\mid X} \).
Theorem 3: Let \( C_n \) be a codebook with \( M \) codewords and composition \( p_X \) with \( p_{\min} \equiv \min_{j : p_X(j) > 0} p_X(j) \). Let \( P_{Y|X} \) be a maximal joint conditional distribution such that and \( P_{Y|X} = W; P_{Y|X} = V \). Let \( \epsilon \geq \frac{2^{H(X)} - 2}{m_{\min}} \) and suppose \( N_{2K+1,p_X}(V) = \{V^1, V^2, \ldots, V^t\} \). Define \( q^i = p_Y^i \cdot V^i \). If for some integer \( a \geq 2 \), for every \( x \in T^n(p_X) \) and for all \( i \in \{1, \ldots, t\} \) we have that

\[
M|T^n_x(V^i)| \geq a^2(n + 1)^{JK - 1} \max_{1 \leq i' \leq t} |T^n(q^{i'})|,
\]

(71)

then, there exists a codeword \( x(m) \in C_n \) such that

\[
\Pr \left[ \hat{m} \neq m \middle| \hat{p}_{y|x(m)} \in N_{2K+1,p_X}(W), x(m) \text{ is sent} \right] > 1 - \frac{2}{a + 1}.
\]

(72)

The above theorem gives us a sphere-packing type of bound. From the method of types we know that \(|T^n_x(V^i)| \geq 2^{nH(V^i|p_X)} \approx 2^{nH(q^i)} \approx 2^{nH(p_X^i)} \cdot V^i \), where the approximation comes from the definition and properties of the neighborhood introduced in Definition 6 (see Section VII for more details). Therefore, inequality (71) approximately implies that

\[
2^{nR}2^{nH(V^i|p_X)} > 2^{nH(p_X^i) V^i},
\]

(73)

or equivalently, \( R > I(p_X, V) \). The theorem states that if approximately \( R > I(p_X, V) \), the error probability of one of the messages is high under \( d \)-decoding.

Proof: From the assumption (71) above we know \( \forall x \in T^n(p_X) \):

\[
M|T^n_x(V^i)| \geq a^2(n + 1)^{JK - 1} \max_{1 \leq i' \leq t} |T^n(q^{i'})|.
\]

(74)

The first step is to show that there is a codeword \( x(m) \in C_n \) such that for all \( 1 \leq i \leq t \) a large proportion of \( T^n_x(V^i) \) have the same conditional type \( V^i \) with at least \( a \) other codewords in \( C_n \), yielding type conflict error. More precisely, we wish to show that there is a codeword \( x(m) \in C_n \) and a family of sets \( F = \{B_i | B_i \in T^n_x(V^i), i = 1, 2, \ldots, t\} \) such that

1) \( |B_i| \geq \frac{a - 1}{a}|T^n_{x(m)}(V^i)| \)

2) \( \forall \hat{y} \in B_i \) there are \( a \) other codewords \( x'(1), x'(2), \ldots, x'(a) \in C_n \) for which \( \hat{p}_{y|x'(1)} = \cdots = \hat{p}_{y|x'(a)} = \hat{p}_{y|x(m)} = V^i \).

This implies that we can find a family of sets \( F = \{B_i\} \) where \( B_i \in T^n_x(V^i) \), such that all members of \( B_i \) for \( 1 \leq i \leq t \) cause type conflict errors.

We prove this result by using proof by contradiction. Suppose there is no such \( x(m) \) with such family \( F = \{B_i\} \). Then, there is no \( x \in C_n \), such that a family \( F = \{B_i\} \) with the above properties exists. Therefore, for any \( x \in C_n \) there is a set \( A_x \) with the following properties:

1) \( A_x \subseteq T^n_x(V^i) \) for some \( 1 \leq i \leq t \),

2) \( |A_x| > \frac{1}{a}|T^n_x(V^i)| \) for the same \( i \) in the above condition,

3) \( \forall \hat{y} \in A_x \) there are at most \( a - 1 \) other codewords \( x'(1), x'(2), \ldots, x'(a - 1) \in C_n \) such that

\[
\hat{p}_{y|x'(1)} = \cdots = \hat{p}_{y|x'(a-1)} = \hat{p}_{y|x}. \]

Now consider any \( 1 \leq i \leq t \) and any \( \hat{y} \in T^n(q^i) \). Using the fact that there are at most \( (n + 1)^{JK - 1} \) conditional types \( p_{Y|X} \) that \( p_Y = q^i \), if this \( \hat{y} \) is a member of more than \( a(n + 1)^{JK - 1} \) sets \( A_x \), we have at
least \( a + 1 \) sets \( A_{\bar{x}(1)} \subseteq T^n_\bar{x}(V^{\bar{i}_1}), \ldots, A_{\bar{x}(a+1)} \subseteq T^n_\bar{x}(V^{\bar{i}_t}) \) for codewords \( \bar{x}(1), \ldots, \bar{x}(a) \in C_n \) and for the same \( 1 \leq i_1 \leq t \). Suppose \( \hat{y} \in A_{\bar{x}(1)} \cap A_{\bar{x}(2)} \cdots \cap A_{\bar{x}(a+1)}. \) Then we have,

\[
\hat{p}_{\hat{y}|\bar{x}(1)} = \hat{p}_{\hat{y}|\bar{x}(2)} = \cdots = \hat{p}_{\hat{y}|\bar{x}(a+1)} = V^{\bar{i}_i}
\]

which contradicts the third condition that the sets \( A_{\bar{x}} \) must satisfy.

As a result, in order not to contradict the third assumption, assume that for all \( 1 \leq i \leq t \) each element of \( T^n(q^i) \) is an element of at most \( a(n + 1)^J(K-1) \) sets \( A_{\bar{x}} \). Now considering the fact that \( A_{\bar{x}} \subset T^n_\bar{x}(V^{\bar{i}}) \) and each element of \( T^n(q^i) \) is in at most \( a(n + 1)^J(K-1) \) sets \( A_{\bar{x}} \) we get the following

\[
\sum_{\bar{x} \in C_n} |A_{\bar{x}}| \leq a(n + 1)^J(K-1) \sum_{i=1}^t |T^n(q^i)| \leq a(n + 1)^J(K-1) \max_{1 \leq i \leq t} |T^n(q^i)|.
\]

On the other hand,

\[
\sum_{\bar{x} \in C_n} |A_{\bar{x}}| > M_a^{-1} |T^n_\bar{x}(V^{\bar{i}})|
\]

where (78) follows from the second property of the sets \( A_{\bar{x}} \) and (79) is from second condition of the Theorem (71). The above two above inequalities give us a contradiction. Therefore, we can find a codeword \( x(m) \in C_n \) with a family of sets \( \mathcal{F} = \{\mathcal{B}_i\} \) such that \( \mathcal{B}_i \subset T^n_\bar{x}(m)(V^{\bar{i}}) \) are large enough and yield type conflict errors with at least 1 codewords.

We proceed by using the assumption that \( P_{Y|X} \) is a maximal joint conditional distribution. Since \( P_{Y|X} \) is maximal, based on Lemma 3 for any \( W \in N_{\bar{x},p_\bar{x}}(W) \) we can find a maximal joint conditional type \( p_{Y|X} \) such that \( p_{Y|X} = W \) and \( p_{Y|X} = V^{\bar{i}} \in N_{\bar{x},p_\bar{x}}(V^{\bar{i}}) \). Construct the graph \( G_\bar{x}(m)(p_{Y|X}) \) for the codeword \( x(m) \) we found above, connecting \( T^n_\bar{x}(m)(W) \) and \( T^n_\bar{x}(m)(V^{\bar{i}}) \). If \( y \in T^n_\bar{x}(m)(W) \) is connected to \( \hat{y} \in T^n_\bar{x}(m)(V^{\bar{i}}) \), by the maximality of \( p_{Y|X} \) and Lemma 3 we have that

\[
d(y, x(m)) \leq d(y, x^{(r)}) \quad \text{for } r = 1, 2, \ldots, a.
\]

The above inequality implies that if \( x(m) \) is transmitted and \( y \in T^n_\bar{x}(m)(W) \) is the channel output, the probability of correct \( d \)-decoding is at most \( \frac{1}{a+1} \) because there are \( a \) other codewords \( x^{(1)}, x^{(2)}, \ldots, x^{(a)} \in C_n \) for which the decoding metric is higher, i.e. \( d(y, x^{(r)}) \leq d(y, x(m)) \) for 1 \( \leq r \leq a \).

Now we count the number of \( y \in T^n_\bar{x}(m)(W) \) that cause a \( d \)-decoding error. Recall that form Definition 4 the set of all \( y \in T^n_\bar{x}(m)(W) \) which are connected to a \( \hat{y} \in \mathcal{B}_i \) in graph \( G_\bar{x}(m)(p_{Y|X}) \) was denoted by \( \Psi_{21}(\mathcal{B}_i) \). In the following, we give a lower bound on \( |\Psi_{21}(\mathcal{B}_i)| \) based on Lemma 1. So far we have proved the following facts:

1) There exists a codeword \( x(m) \in C_n \) and a family of sets \( \mathcal{F} = \{\mathcal{B}_i\} \) such that \( \mathcal{B}_i \subset T^n_\bar{x}(m)(V^{\bar{i}}) \) and \( |\mathcal{B}_i| \geq \frac{a}{a+1} |T^n_\bar{x}(m)(V^{\bar{i}})| \).

2) \( \forall \hat{y} \in \mathcal{B}_i \) connected to \( y \in T^n_\bar{x}(m)(W) \) in graph \( G_\bar{x}(m)(p_{Y|X}) \) the following holds

\[
\mathbb{P}[\hat{m} \neq m \mid y \text{ is received, } x(m) \text{ is sent}] \geq \frac{a}{a+1}.
\]
Now we count the number of elements of $\Psi_{21}(B_i)$ in $G_{x(m)}(p_{Y|X})$ for $1 \leq i \leq t$, since the $d$-decoder makes errors on elements of $\Psi_{21}(B_i)$. Now using the fact that $|B_i| \geq \frac{a - 1}{a} |T_x^n(V')|$ and using Lemma 1 with $\mathcal{V}_1 = T_x^n(W)$ and $\mathcal{V}_2 = T_x^n(V')$ we get

$$\frac{|\Psi_{21}(B_i)|}{|T_x^n(W)|} \geq \frac{|B_i|}{|T_x^n(V')|}$$

(82)

$$\geq \frac{a - 1}{a} \frac{|T_x^n(V')|}{|T_x^n(W)|}$$

(83)

$$= \frac{a - 1}{a}.$$  

(84)

In the remaining part of the proof we relate $\frac{|\Psi_{21}(B_i)|}{|T_x^n(W)|}$ to the probability of error. Suppose $x(m)$ is sent over the channel and $y$ is received. Now note by the definition of conditional type, all elements of $T_x^n(W)$ are equally likely to appear at the output of the channel when $x(m)$ is sent. Therefore, for every $y_0 \in T_x^n(W)$,

$$P[y \in T_x^n(W) | x(m) \text{ is sent}] = \sum_{y \in T_x^n(W)} W^n(y|x(m))$$

(85)

$$= |T_x^n(W)| \cdot W^n(y_0|x(m))$$

(86)

where (85) follows since $W^n(y|x(m))$ is the same for all $y \in T_x^n(W)$. Now we get the following inequality

$$P[\hat{m} \neq m | \hat{p}_{y|x(m)} = W, x(m) \text{ is sent}] = P[\hat{m} \neq m | y \in T_x^n(W), x(m) \text{ is sent}]$$

(87)

$$= \frac{P[\hat{m} \neq m, y \in T_x^n(W) | x(m) \text{ is sent}]}{P[y \in T_x^n(W) | x(m) \text{ is sent}]}$$

(88)

$$= \frac{\sum_{y \in T_x^n(W)} P[\hat{m} \neq m, y \neq \hat{y} | x(m) \text{ is sent}]}{P[y \in T_x^n(W) | x(m) \text{ is sent}]}$$

(89)

$$= \frac{\sum_{y \in T_x^n(W)} P[\hat{m} \neq m | y = \hat{y}, x(m) \text{ is sent}] \cdot W^n(\hat{y}|x(m))}{P[y \in T_x^n(W) | x(m) \text{ is sent}]}$$

(90)

$$= \frac{1}{|T_x^n(W)|} \sum_{y \in T_x^n(W)} P[\hat{m} \neq m | y = \hat{y}, x(m) \text{ is sent}]$$

(91)

$$\geq \frac{1}{|T_x^n(W)|} \sum_{y \in \Psi_{21}(B_i)} P[\hat{m} \neq m | y = \hat{y}, x(m) \text{ is sent}]$$

(92)

$$\geq \frac{|\Psi_{21}(B_i)|}{|T_x^n(W)|} \cdot \frac{a}{a + 1}$$

(93)

$$\geq \frac{a - 1}{a + 1}$$

(94)

where (88) follows from definition of conditional probability, (91) follows from (86), (92) follows since $\Psi_{21}(B_i) \subseteq T_x^n(W)$, (93) follows from (81) and (94) from (84).

In the last step we have the following inequality,

$$P[\hat{m} \neq m | \hat{p}_{y|x(m)} \in N_{p_x}(W), x(m) \text{ is sent}]$$

$$= \frac{P[\hat{m} \neq m, \hat{p}_{y|x(m)} \in N_{p_x}(W) | x(m) \text{ is sent}]}{P[\hat{p}_{y|x(m)} \in N_{p_x}(W) | x(m) \text{ is sent}]}$$

(95)
where (95) follows from the definition of conditional probability, (98) follows from inequality (94) and (99) follows from the definition of \( \mathcal{N}_{Y}^{W} (W) \). This concludes the proof.

\[ \]

VI. FROM TYPES TO DISTRIBUTIONS

It is known that if rate \( R > 0 \) is achievable then for any \( \epsilon > 0 \) there exists constant composition codes of rate \( R - \epsilon \) whose probability of error tends to 0. In the following lemma, we prove that if rate \( R \) is achievable, then, for any \( \epsilon > 0 \) there exist constant composition codes of rate \( R - \epsilon \) with vanishing probability of error that have the additional property that their composition \( p_{n} \in \mathcal{P}_{n}(\mathcal{X}) \) is such that if \( p_{n}(j) > 0 \), then \( p_{n}(j) \geq \delta \) for \( \delta > 0 \) independent of \( n \), for all \( j = 1, \ldots, J \).

**Definition 7:** Let \( \tilde{C}_{n} \) be a codebook. We say that \( \tilde{C}_{n} \) is a \( \delta \)-reduction of \( C_{n} \) if there exists a sub-codebook \( \tilde{C}_{n} \subseteq C_{n} \) of composition \( p_{X} \in \mathcal{P}_{n}(\mathcal{X}) \) such that \( \tilde{C}_{n} \) is obtained by eliminating all symbols in the set \( \mathcal{I} = \{ j \in \mathcal{X} | p_{X}(j) < \delta \} \) from \( \tilde{C}_{n} \).

**Lemma 5:** Let \( R > 0 \) be a rate, then for any \( \epsilon > 0 \) there exists a \( \delta > 0 \) independent of \( n \) such that for any codebook \( C_{n} \) of rate \( R \) there exists a \( \delta \)-reduction constant composition codebook \( \tilde{C}_{n} \) with the following properties:

\[ \hat{n} \geq (1 - (|\mathcal{X}| - 1)\delta)n \]  
\[ P_{e,\text{max}}^{d}(\tilde{C}_{n}) \leq P_{e,\text{max}}^{d}(C_{n}) \]  
\[ \frac{1}{n} \log(|\hat{C}_{n}|) \geq \frac{1}{n} \log(|C_{n}|) - \epsilon + O \left( \frac{\log(n)}{n} \right). \]

**Proof:** For any \( n > 0 \) we know \(|\mathcal{P}^{n}(\mathcal{X})| \leq (n + 1)^{|\mathcal{X}| - 1} \). Therefore, by the pigeonhole principle, any codebook \( C_{n} \) contains a constant composition sub-codebook \( \tilde{C}_{n} \) of type \( p_{n} \) and such that \( |\tilde{C}_{n}| \geq \frac{c_{n}}{(n + 1)^{|\mathcal{X}| - 1}} \) codewords. Let \( \mathcal{I} = \{ i_{1}, i_{2}, \ldots, i_{t} \} \subseteq \mathcal{X} \) be the set of all symbols \( j \in \mathcal{X} \) that \( p_{n}(j) < \delta \). Then, there are

\[ \binom{n}{n p_{n}(i_{1}), n p_{n}(i_{2}), \ldots , n p_{n}(i_{t})} = \binom{n}{n p_{n}(i_{1})!(n p_{n}(i_{2})!) \cdot \cdots \cdot (n p_{n}(i_{t})!)(n - \sum_{j=1}^{t} n p_{n}(i_{j}))!} \]

different places for symbols of set \( \mathcal{I} \) in a string of length \( n \). For ease of notation we use the following notation

\[ \binom{n}{n p_{n}(\mathcal{I})} = \binom{n}{n p_{n}(i_{1}), n p_{n}(i_{2}), \ldots , n p_{n}(i_{t})}. \]

As a result, by again using the pigeonhole principle, there exists a sub-codebook \( \tilde{C}_{n} \subseteq C_{n} \) with \( |\tilde{C}_{n}| \geq \frac{c_{n}}{(n + 1)^{|\mathcal{X}| - 1}} \) codewords where all symbols in set \( \mathcal{I} \) are in the same position. By being in the same position we mean that the codewords of \( \tilde{C}_{n} \) have all symbols \( i_{1}, i_{2}, \ldots , i_{t} \) in the same position. Let \( Z \subseteq \{ 1, \ldots , n \} \)
be set of positions where symbols in $I$ are placed. We then form the $\delta$-reduced codebook $\hat{C}_n$ by shortening the codewords of $C_n$ such that symbols in positions in $Z$ are removed. The rate of this codebook is therefore

$$\frac{1}{n} \log(|\hat{C}_n|) \geq \frac{1}{n} \log(|C_n|) \geq \frac{1}{n} \log \left( \frac{|C_n|}{(n+1)^{|X|}-1(n \cdot p_n(Z))} \right)$$

(105)

By using Stirling’s factorial formula we upper-bound the contribution of the multinomial coefficient by the entropy as follows

$$\frac{1}{n} \log \left( \frac{n}{n \cdot p_n(I)} \right) = \frac{1}{n} \log \left( \frac{n!}{(np_n(i_1))! \cdots (np_n(i_t))!} \right)$$

$$\leq H(p_n(i_1), p_n(i_2), \cdots, p_n(i_t), 1 - \sum_{j=1}^t p_n(i_j)) + O \left( \frac{\log(n)}{n} \right)$$

(106)

$$\leq H(\delta, \delta, \cdots, \delta, 1 - (|X| - 1)\delta) + O \left( \frac{\log(n)}{n} \right)$$

(108)

where $H(q_1, \cdots, q_m) = -\sum_{i=1}^m q_i \log q_i$ denotes the entropy function of probability mass function with $m$ nonzero mass points with probabilities $q_1, \cdots, q_m$ and (108) follows from observing that $p_n(i_1), p_n(i_2), \cdots, p_n(i_t) \leq \delta,$ $t \leq |X| - 1$ and the fact that $\delta$ can be chosen sufficiently small.

Summarizing, we get the following inequality

$$\frac{1}{n} \log(|\hat{C}_n|) \geq \frac{1}{n} \log(|C_n|) \geq R - H(\delta, \delta, \cdots, \delta, 1 - (|X| - 1)\delta) - (|X| - 1) \frac{\log(n + 1)}{n} + O \left( \frac{\log(n)}{n} \right)$$

(109)

(110)

Now choosing $\delta$ in a way that $H(\delta, \delta, \cdots, \delta, 1 - (|X| - 1)\delta) < \varepsilon$ we get the desired result.

It remains is to show that

$$P_{e, \text{max}}^d(\hat{C}_n) \leq P_{e, \text{max}}^d(C_n).$$

(111)

This directly follows from the fact that all symbols in $I$ are in the same position in the codebook $\hat{C}_n$. Let us define $\alpha : \mathcal{Y}^n \rightarrow \hat{\mathcal{Y}}^n$ as the function that gets a string $y \in \mathcal{Y}^n$ and gives $\alpha(y) \in \hat{\mathcal{Y}}^n$ by eliminating the symbols in positions in the set $Z$. Moreover, let $\mu : \hat{C}_n \rightarrow \hat{C}_n$ that performs the same operations on the codewords of $\hat{C}_n$. Then, for any $y \in \hat{\mathcal{Y}}^n$ and $x \in \hat{C}_n$ we have $d(x, y) - d(\mu(x), \alpha(y))$ is a function of $y$, because all the codewords in $\hat{C}_n$ have the same symbols at the eliminated entries. As a result, if $y$ is decoded to $x(\hat{m}) \in \hat{C}_n$ under $d$-decoding, then $\alpha(y)$ would be decoded to $\mu(x(\hat{m}))$ under $d$-decoding. Let $x_Z$ and $y_Z$ be the symbols of $x$ and $y$ in positions in set $Z$, respectively. Now notice that for all $y' \in \hat{\mathcal{Y}}^n$ and all $x \in \hat{C}_n$

$$\sum_{y' \in \alpha^{-1}(y')} W^n(y|x) = \sum_{y_Z} W^{\hat{n}}(y'|\mu(x)) W^{n-\hat{n}}(y_Z|x_Z)$$

$$= W^{\hat{n}}(y'|\mu(x)) \sum_{y_Z} W^{n-\hat{n}}(y_Z|x_Z)$$

(112)

(113)

$$= W^{\hat{n}}(y'|\mu(x))$$

(114)
As a result we have

\[ W_1 = \varepsilon \frac{1}{n} X \]

the random variable symbols equal to \( j \) for the output of channel \( W \) when \( x \) is sent. Then, \( \forall \gamma > 0 \) we have

\[
P[|\hat{X} - E[\hat{X}]| \geq \gamma] \leq e^{-2n\gamma^2} \tag{117}\]

The above result is proved because in order to use the following theorem, we need the frequency of each symbol in any codeword to be proportional to \( n \).

**Theorem 4 (Hoeffding’s inequality):** Assume \( X_1, i = 1, 2, \cdots, n \) are independent random variables taking values on \([0, 1]\). Let \( \bar{X} = \frac{1}{n}(X_1 + X_2 + \cdots + X_n) \). Then \( \forall \gamma > 0 \)

\[ P[|X - E[X]| \geq \gamma] \leq e^{-2n\gamma^2} \tag{118} \]

The following lemma is the statement that we need to show that empirical conditional type of the received sequence given the sent message would be close to \( W \).

**Lemma 6:** Let \( x \in T^n(p_X) \) be a codeword, and denote by \( y \) the output of channel \( W \) when \( x \) is sent. Then, \( \forall \gamma > 0 \) we have

\[
P[\hat{p}_{y|x}(j,k) \in \mathcal{N}_{\gamma,p_X}(W) \mid x \text{ is sent}] > 1 - |\mathcal{X}||\mathcal{Y}| \cdot e^{-2np_{\text{min}}\gamma^2}. \tag{119}\]

**Proof:** Let \( (j,k) \in \mathcal{X} \times \mathcal{Y} \) and assume \( p_x(j) > 0 \). We know from the definition of types there are \( np_X(j) \) symbols equal to \( j \) in \( x \). Without loss of generality assume, \( x_1 = x_2 = \cdots = x_{np_X(j)} = j \). Now define the random variable \( X_i, i = 1, 2, \cdots, np_X(j) \) in the following way

\[
X_i = \begin{cases} 1 & (y_i, x_i) = (k, j) \\ 0 & \text{otherwise.} \end{cases} \tag{120}\]

As a result, the conditions of Hoeffding’s inequality hold for \( X_i, i = 1, 2, \cdots, np_X(j) \) and \( E[X_i] = P[X_i = 1] = W(k|j) \). Therefore, we get the following

\[
P[|W(j,k) - \hat{p}_{y|x}(j,k)| \geq \gamma \mid x \text{ is sent}] \leq e^{-2np_X(j)\gamma^2} \tag{121}\]

Now, lower bounding \( p_X(j) \) by \( p_{\text{min}} \) we get

\[
P[|W(j,k) - \hat{p}_{y|x}(j,k)| \geq \gamma \mid x \text{ is sent}] \leq e^{-2np_{\text{min}}\gamma^2}. \tag{122}\]

As a result we have

\[
P[\hat{p}_{y|x}(j,k) \in \mathcal{N}_{\gamma,p_X}(W) \mid x \text{ is sent}] = 1 - P[\bigcup_{j,p_x(j)>0,k} \{|W(j,k) - \hat{p}_{y|x}(j,k)| > \gamma \} \mid x \text{ is sent}] \geq 1 - \sum_{j,p_x(j)>0,k} P[|W(j,k) - \hat{p}_{y|x}(j,k)| > \gamma \mid x \text{ is sent}] \geq 1 - \sum_{j,p_x(j)>0,k} e^{-2np_X(j)\gamma^2} \geq 1 - |\mathcal{X}||\mathcal{Y}|e^{-2np_{\text{min}}\gamma^2}, \tag{123}\]

where (115) follows from the fact that \( W \) is a memoryless channel. Moreover, \( x \in \mathcal{Z} \) and \( y \) are strings consisting of symbols of the index set \( \mathcal{X} \) of \( x \) and \( y \) respectively. As a result, the probability of error of any codeword \( x \in \mathcal{C}_n \) is equal to probability of error of \( \mu(x) \in \mathcal{C}_n \). Thus,

\[
P_{e,\text{max}}(\mathcal{C}_n) = P_{e,\text{max}}(\hat{\mathcal{C}}_n) \tag{115}\]

Since \( \mathcal{C}_n \) is a sub-codebook of \( \hat{\mathcal{C}}_n \),

\[
P_{e,\text{max}}(\mathcal{C}_n) \leq P_{e,\text{max}}(\hat{\mathcal{C}}_n). \tag{116}\]

Combining (115) and (116) completes the proof. 

\[ \square \]
where (123) follows from (121).

The above result shows that when the frequency of every symbol in the codebook grows proportional to \( n \), then conditional type of the output string given the sent message will be close to \( W \) with high probability.

### VII. Proof of the Main Theorem

In this part we prove the final part of Theorem 1 using the material developed in the previous sections. Assume \( R = \bar{R}_d(W) + \sigma \) for some \( \sigma > 0 \). Now, choose \( \epsilon > 0 \) small enough such that if \( |V - V|_\infty \leq 2K\epsilon \) for conditional distribution \( V, V \), then for any distribution \( p \) on \( X \) we have that

\[
|H(V|p) - H(V|p)| < \frac{\sigma}{4},
\]

(126)

\[
|H(p^T V') - H(p^T V)| < \frac{\sigma}{4}.
\]

(127)

The reader is referred to the Appendix for a discussion on the choice of \( \epsilon \).

From Lemma 5 with \( \epsilon = \frac{\sigma}{4} \), for any codebook \( C_n \) with \( M \geq 2^nR \) codewords, there exists a \( \delta \)-reduction constant composition codebook \( \hat{C}_n \) of length \( \tilde{n} \) and type \( \hat{p}_n \) such that (100)–(102) are satisfied. Since the required \( \delta \) to satisfy the above inequalities is independent of \( n \), then choose \( N_0 \) large enough such that \( \epsilon \geq \frac{2|Y|}{n_0(1-|X|^{-1})M} \).

Set \( n > N_0 \). Now choose a maximal joint conditional distribution \( P_{Y|X} \) such that \( I(\hat{p}_n, P_{Y|X}) \leq \bar{R}_d(W) \) and let \( V = P_{Y|X} \). Such a \( P_{Y|X} \) exists because the set \( \mathcal{M}_{\text{max}}(d) \cap \{ P_{Y|X} = W \} \) which is the domain of the minimization in (17) is a compact set and the minimizer always exists. Now, for any conditional distributions \( \hat{V} \) such that \( |V' - V|_\infty \leq 2K\epsilon \)

\[
\max_{\nabla \in \mathcal{N}_{2K^2, \hat{p}_n}(V)} H(p^T \nabla) - H(\hat{V}|\hat{p}_n) \leq |H(p^T V) - H(V|\hat{p}_n)| + \frac{\sigma}{2}
\]

(128)

\[
= I(\hat{p}_n, V) + \frac{\sigma}{2}
\]

(129)

\[
= I(\hat{p}_n, V) + \frac{\sigma}{2}
\]

(130)

where (129) follows from (126) and (127).

Now suppose \( \mathcal{N}_{\epsilon, \hat{p}_n}(V) = \{ V^1, V^2, \cdots, V^t \} \) and \( q^i = \hat{p}_n V^i. \) Now for any \( 1 \leq i \leq t \)

\[
\frac{1}{n} \log \max_{1 \leq s \leq t} \frac{|T^n_s(q^i)|}{|T^n_s(V)|} = \frac{1}{n} \log \frac{2^n H(q^i) + O(\frac{\log(n)}{n})}{2^n H(V) + O(\frac{\log(n)}{n})}
\]

(131)

\[
\leq I(\hat{p}_n, V) + \frac{\sigma}{2} + O\left( \frac{\log(n)}{n} \right)
\]

(132)

where \( i' = \arg \max_{1 \leq s \leq t} |T^n_s(q^i)| \) and (132) follows form (130). Now, for \( n > N_0 \) we have from (102) with \( \epsilon = \frac{\sigma}{4} \) and the condition \( I(\hat{p}_n, P_{Y|X}) \leq \bar{R}_d(W) \) that

\[
\frac{|T^n_s(V)|}{\max_{1 \leq s \leq t} |T^n_s(q^i)|} \geq 2^n(R - \frac{s}{2} - I(\hat{p}_n, V) - \frac{\sigma}{2} + O(\frac{\log(n)}{n}))
\]

(133)

As a result,

\[
|\hat{C}_n||T^n_s(V)| \geq 2^n(R - \frac{s}{2} + O(\frac{\log(n)}{n})) \max_{1 \leq s \leq t} |T^n_s(q^i)|.
\]

(134)

Setting \( a = \left[ \frac{2^n(R - \frac{s}{2} + O(\frac{\log(n)}{n}))}{(n+1)2^n(R - \frac{s}{2})} \right] \) validates the conditions of Theorem 3. As a result, there exists \( x(m) \in \hat{C}_n \) such that

\[
P[\hat{m} \neq m | \hat{p}_{g|x(m)} \in \mathcal{N}_{\epsilon, \hat{p}_n}(W), x(m) \text{ is sent}] > 1 - \frac{2}{a+1}.
\]

(135)
Choosing $N_1$ such that if $n > N_1$ is large enough, we bound
\[
\lambda \geq \frac{1}{2} \cdot \frac{2^n (\hat{\lambda} + O(\frac{\log(n)}{n}))}{(n + 1)(JK - 1)} \geq 2^n (\hat{\lambda} + O(\frac{\log(n)}{n})) - (JK - 1) \frac{1}{2n} \frac{\log(n)}{n} \geq 1
\] (136)
\[
\lambda \geq \frac{1}{2n} \frac{\log(n)}{n} \geq 1
\] (137)
Finally, we write,
\[
P_{e,\text{max}}(C_n) \geq P_{e,\text{max}}(\tilde{C}_n)
\] (138)
\[
= \max_{m \in \{1, \ldots, M\}} \mathbb{P} \left[ \hat{m} \neq m \mid \mathbf{x}(m) \text{ is sent} \right]
\] (139)
\[
\geq \max_{m \in \{1, \ldots, M\}} \mathbb{P} \left[ \hat{m} \neq m \mid \hat{P}_{y|x(m)} \in \mathcal{N}_{\tilde{Z}} \cdot P_k(W), \mathbf{x}(m) \text{ is sent} \right] \cdot \mathbb{P} \left[ \hat{P}_{y|x(m)} \in \mathcal{N}_{\tilde{Z}} \cdot P_k(W) \mid \mathbf{x}(m) \text{ is sent} \right]
\] (140)
\[
\geq \left( 1 - \frac{2\lambda}{a + 1} \right) \left( 1 - |X||Y|2^{-2n\lambda^2 a^2} \right)
\] (141)
\[
\geq 1 - 2^{-\hat{\lambda} L_{sc}(R)}
\] (142)
where (141) follows from applying Theorem 3 to the first probability in (140) and Lemma 6 to the second probability in (140), where $P_{e,\text{sc}}(R) \triangleq \min \left\{ \frac{\delta^2}{\lambda^2} - \log \frac{|X||Y|}{n}, \frac{1}{2} \left( 1 + O \left( \frac{\log(n)}{n} \right) \right) \right\}$. Setting $\lambda$ larger than max\{$N_0, N_1$\} yields the desired result.

VIII. CONVEXITY ANALYSIS

In this section, we show that the optimization problem (17) is a convex-concave saddlepoint problem. First we argue that the constraints induce a convex set.

**Lemma 7:** For any channel $W$ and metric $d$, the set of joint conditional distributions $P_{Y|X}$ satisfying both $P_{Y|X} \in \mathcal{M}_{\text{max}}(d)$ and $P_{Y|X} = W$, is a convex set.

**Proof:** Let $P_{Y|X}$ and $P'_{Y|X}$ both satisfy the above constraints. Now for any $0 < \lambda < 1$ we have,
\[
\lambda P_{Y|X} + (1 - \lambda) P'_{Y|X} = W
\] (143)
In addition, if for some $k_1, k_2$ we have $j \notin S_2(k_1, k_2)$, both $P_{Y|X}(k_1, k_2|j)$ and $P'_{Y|X}(k_1, k_2|j)$ are equal to zero, and so is any linear combination of them. Therefore,
\[
\lambda P_{Y|X} + (1 - \lambda) P'_{Y|X} \in \mathcal{M}_{\text{max}}(d)
\] (144)
Moreover, $I(P_X, P_{Y|X})$ is convex in terms of $P_{Y|X}$, and concave in terms of $P_X$. Since $P_{Y|X}$ is a linear function of $P_{Y|X}$, we get that $I(P_X, P_{Y|X})$ is also convex in terms of $P_{Y|X}$. Therefore from the minimax theorem [15],
\[
\hat{R}_d(W) = \max_{P_X} \min_{P_{Y|X} \in \mathcal{M}_{\text{max}}(d)} I(P_X, P_{Y|X})
\] (145)
\[
= \min_{P_{Y|X} \in \mathcal{M}_{\text{max}}(d)} \max_{P_X} I(P_X, P_{Y|X})
\] (146)
\[
= \min_{P_{Y|X} \in \mathcal{M}_{\text{max}}(d)} C(P_{Y|X}).
\] (147)
The rest of this section is devoted to deriving the KKT conditions for the optimization problem in (17). Given that \( I(P_X, P_{Y|X}) \) is convex in \( P_{Y|X} \) and concave in \( P_X \), then the KKT conditions are sufficient for global optimality (16). For convenience, we define \( f(P_X, P_{Y|X}) \triangleq I(P_X, P_{Y|X}) \) and rewrite the optimization problem in (17) as,

\[
R_d(W) = \max_{P_X} \min_{P_{Y|X} \in \mathcal{M}_{\text{max}}(d)} f(P_X, P_{Y|X}).
\]

(148)

Let \( P_X^*, P_{Y|X}^* \) be the optimal input and joint conditional distributions in (148) and \( q_Y^* \) be the output distribution induced by \( P_X^* \) and \( P_{Y|X}^* \). Then for \( P_X^* \) we have the following constraints:

\[
P_X^*(j) \geq 0, \quad \forall j \in \mathcal{X}
\]

(149)

\[
\sum_{j \in \mathcal{X}} P_X^*(j) = 1.
\]

(150)

Let \( \mu_j, j = 1, 2, \cdots, J \) be the Lagrange multipliers corresponding the inequalities in (149) and \( \rho \) be the Lagrange multiplier corresponding to (150). Therefore, from stationarity we have

\[
\frac{\partial}{\partial P_X(j)} f(P_X, P_{Y|X}^*) \bigg|_{P_X = P_X^*} = \rho + \mu_j
\]

(151)

Then from the complementary slackness we have \( \mu_j P_X^*(j) = 0 \) and from the dual feasibility we have \( \mu_j \geq 0 \) which leads to the separation of the equations (151) in two cases (16). If \( P_X^*(j) > 0 \)

\[
\frac{\partial}{\partial P_X(j)} f(P_X, P_{Y|X}^*) \bigg|_{P_X = P_X^*} = \rho
\]

(152)

And when \( P_X^*(j) = 0 \) we have

\[
\frac{\partial}{\partial P_X(j)} f(P_X, P_{Y|X}^*) \bigg|_{P_X = P_X^*} \leq \rho
\]

(153)

Note that, because there is no other constraint on \( \mu_j \), all of the KKT conditions are summarized in (152) and (153). Moreover, computing the derivatives in (152) and (153) gives

\[
\frac{\partial}{\partial P_X(j)} f(P_X, P_{Y|X}^*) \bigg|_{P_X = P_X^*} = \sum_{k \in \mathcal{Y}} P_{Y|X}^*(k|j) \log \frac{P_{Y|X}^*(k|j)}{q_Y^*(k)} - 1.
\]

(154)

Similarly, for \( P_{Y|X}^* \) we have the following constraints. For all \( j, k_1, k_2 \in \mathcal{X} \times \mathcal{Y} \times \mathcal{Y} \)

\[
P_{Y|X}^*(k_1, k_2|j) \geq 0,
\]

(155)

\[
P_{Y|X}^*(k_1, k_2|j) = 0, \quad \text{if } j \notin \mathcal{S}_d(k_1, k_2)
\]

(156)

where (155) corresponds to \( P_{Y|X}^*(k_1, k_2|j) \) being a distribution and (156) corresponds to \( P_{Y|X}^*(k_1, k_2|j) \in \mathcal{M}_{\text{max}}(d) \). Moreover from the constraint \( P_{Y|X} = W \) we get for all \( j, k_1 \in \mathcal{X} \times \mathcal{Y} \)

\[
\sum_{k_2} P_{Y|X}^*(k_1, k_2|j) = W(k_1|j)
\]

(157)

For the ease of notation, we skip the step of explicitly considering a Lagrange multiplier for (155). Details follow similarly to the above derivation. If we use a Lagrange multiplier \( \lambda_{j,k_1} \) for each of the conditions in (157), we have when \( P_{Y|X}^*(k_1, k_2|j) > 0 \)

\[
\frac{\partial}{\partial P_{Y|X}^*(k_1, k_2|j)} f(P_X^*, P_{Y|X}^*) \bigg|_{P_{Y|X} = P_{Y|X}^*} = \lambda_{j,k_1}
\]

(158)
and when \( P^*_Y|X(k_1, k_2|j) = 0 \) and \( j \in S_d(k_1, k_2) \) we should have

\[
\frac{\partial}{\partial P^*_Y|X(k_1, k_2|j)} f(P^*_X, P^*_Y|X) \bigg|_{P^*_Y|X = P^*_Y|X} \geq \lambda_{j,k_1}.
\] (159)

Explicitly computing the derivative gives

\[
\frac{\partial}{\partial P^*_Y|X(k_1, k_2|j)} f(P^*_X, P^*_Y|X) \bigg|_{P^*_Y|X = P^*_Y|X} = P^*_X(j) \log \frac{P^*_Y|X(k_2|j)}{q^*_Y(k_2)}. \] (160)

Summarizing, for the KKT optimality conditions of we get the following inequalities

1) If \( P^*_X(j) > 0 \),

\[
\sum_{k \in Y} P^*_Y|X(k|j) \log \frac{P^*_Y|X(k|j)}{q^*_Y(k)} = 1 + \rho
\] (161)

2) If \( P^*_X(j) = 0 \),

\[
\sum_{k \in Y} P^*_Y|X(k|j) \log \frac{P^*_Y|X(k|j)}{q^*_Y(k)} \leq 1 + \rho
\] (162)

3) If \( P^*_Y|X(k_1, k_2|j) > 0 \),

\[
P^*_X(j) \log \frac{P^*_Y|X(k_2|j)}{q^*_Y(k_2)} = \lambda_{j,k_1}
\] (163)

4) If \( P^*_Y|X(k_1, k_2|j) = 0 \) and \( j \in S_d(k_1, k_2) \),

\[
P^*_X(j) \log \frac{P^*_Y|X(k_2|j)}{q^*_Y(k_2)} \geq \lambda_{j,k_1}
\] (164)

In the next sections, we employ the above KKT conditions to efficiently compute \( \tilde{R}_d(W) \) and to analyze the multiletter version of the bound.

### IX. Computation of \( \tilde{R}_d(W) \)

In this section, we turn to the computation of the proposed upper bound \( \tilde{R}_d(W) \). Before describing the algorithm in detail, we introduce a number of concepts related to convex-concave optimization. Let \( D \subset \mathbb{R}^n \) be an open convex set. The standard inner product on \( \mathbb{R}^n \) is denoted by \( \langle \cdot, \cdot \rangle \). A mirror map is a function \( \Psi : D \to \mathbb{R} \) with the following properties:

1) \( \Psi \) is strictly convex and continuously differentiable on \( D \), where strict convexity means that for all \( v_1, v_2 \in D \)

\[
\Psi(v_1) - \Psi(v_2) - \langle \nabla \Psi(v_2), v_1 - v_2 \rangle > 0,
\] (165)

2) The range of \( \nabla \Psi \) is all of \( \mathbb{R}^n \) i.e. \( \nabla \Psi(D) = \mathbb{R}^n \),

3) The gradient of \( \Psi \) diverges on the boundary of \( D \), denoted by \( \partial D \), that is

\[
\lim_{w \to \partial D} \| \nabla \Psi(v) \| = \infty,
\] (166)

where \( \| : \| \) denotes the Euclidean norm.

The Bregmann divergence \( B_\Psi(\cdot, \cdot) : D \times D \to \mathbb{R} \) with respect to a mirror map \( \Psi \) is defined as

\[
B_\Psi(v_1, v_2) = \Psi(v_1) - \Psi(v_2) - \langle \nabla \Psi(v_2), v_1 - v_2 \rangle.
\] (167)
Let $\mathcal{D} \subseteq \mathbb{R}^n$ be a convex set. Function $h : \mathcal{D} \to \mathbb{R}$ is said to be $\alpha$-strongly convex with respect to norm $| \cdot |$ if it is differentiable on $\mathcal{D}$ and for all $v_1, v_2 \in \mathcal{D}$ we have

$$h(v_1) - h(v_2) - \langle \nabla h(v_2), v_1 - v_2 \rangle \geq \frac{\alpha}{2} |v_1 - v_2|^2,$$

(168)

where the norm $| \cdot |$ is not necessarily induced by the standard inner product, i.e. it is not necessarily the Euclidean norm. If the mirror map $\Psi : \mathcal{D} \to \mathbb{R}$ is 1-strongly convex w.r.t. then from the definition (168) for all $v_1, v_2 \in \mathcal{D}$ we have

$$B_\Psi(v_1, v_2) \geq \frac{1}{2} |v_1 - v_2|^2.$$

(169)

We aim to compute the value of the following saddlepoint problem.

$$\hat{R}_d(W) = \max_{P_X} \min_{P_{Y \mid X} \in \mathcal{M}_{\text{max}}(d)} f(P_X, P_{Y \mid X}).$$

(170)

For ease of notation and consistency we define $\mathcal{E}_1$ and $\mathcal{E}_2$ be the constraint sets corresponding to the maximization and minimization, respectively,

$$\mathcal{E}_1 = \left\{ v \in \mathbb{R}^J \mid v(j) \geq 0, \sum_{j=1}^J v(j) = 1 \right\},$$

(171)

$$\mathcal{E}_2 = \left\{ u \in \mathbb{R}^{J \times K \times K} \mid \sum_{k_2=1}^K u(j, k_1, k_2) = W(k_1 \mid j), u(j, k_1, k_2) \geq 0, u(j, k_1, k_2) = 0 \text{ if } j \notin \mathcal{S}_d(k_1, k_2) \right\},$$

(172)

where $\mathcal{E}_1$ corresponds to the set of distributions over $\mathcal{X}$ and $\mathcal{E}_2$ corresponds to the set of maximal joint conditional distributions $P_{Y \mid X}$ such that $P_{Y \mid X} = W$ i.e. $\mathcal{M}_{\text{max}}(d) \cap \{ P_{Y \mid X} \mid P_{Y \mid X} = W \}$. There is a natural bijection among the two sets by corresponding $u$ to $P_{Y \mid X}$ whenever for all $j, k_1, k_2 \in \mathcal{X} \times \mathcal{Y} \times \mathcal{Y}$ we have $u(j, k_1, k_2) = P_{Y \mid X}(k_1, k_2 \mid j)$. With a slight abuse of notation let $f$ be defined for vectors $v \in \mathcal{E}_1, u \in \mathcal{E}_2$ as it is defined for their corresponding distributions $P_X, P_{Y \mid X}$ in the previous section i.e. $f(v, u) \triangleq f(P_X, P_{Y \mid X})$. Therefore, with a slight abuse of notation, we rewrite the saddlepoint problem (170) as

$$\hat{R}_d(W) = \max_{v \in \mathcal{E}_1} \min_{u \in \mathcal{E}_2} f(v, u).$$

(173)

In the rest of this section whenever $u$ is used it is considered that $u(j, k_1, k_2) = 0$ if $j \notin \mathcal{S}_d(k_1, k_2)$, i.e., that the corresponding $P_{Y \mid X} \in \mathcal{M}_{\text{max}}(d)$. Additionally, we choose

$$\mathcal{D}_1 = \{ v \in \mathbb{R}^J \mid 0 \leq v(j), 0 \leq j \leq J \}$$

(174)

and

$$\mathcal{D}_2 = \{ u \in \mathbb{R}^{J \times K \times K} \mid 0 \leq u(j, k_1, k_2), 0 \leq j \leq J, 0 \leq k_1, k_2 \leq K, u(j, k_1, k_2) = 0 \text{ if } j \notin \mathcal{S}_d(k_1, k_2) \}. $$

(175)

It is known that the function $\Psi_1(v) = \sum_i v(i) \log v(i)$ is 1-strongly convex mirror maps on $\mathcal{D}_1$ w.r.t. norm $| \cdot |_1$ [17]. Additionally, let $\Psi_2(u) = \sum_{j, k_1, k_2} I(j \in \mathcal{S}_d(k_1, k_2))u(j, k_1, k_2) \log v(j, k_1, k_2)$. Note that $\nabla \Psi_2$ is surjective on $\{ u \in \mathbb{R}^{J \times K \times K} \mid u(j, k_1, k_2) = 0 \text{ if } j \notin \mathcal{S}_d(k_1, k_2) \}$. Moreover, in all of the computations regarding $u$ we only use the entries $u(j, k_1, k_2)$ such that $j \in \mathcal{S}_d(k_1, k_2)$ and ignore all other entries, i.e. they are set to 0 from the beginning of the algorithm and never change. Therefore, with a slight abuse of notation we say $\Psi_2$ is 1-strongly convex mirror map on $\mathcal{D}_2$ w.r.t. norm $| \cdot |_1$. Note that for $\Psi_2$ being a mirror map, from
Note that when \( \gamma \in \mathcal{S}_d(k_1, k_2) \) and \( \nabla \Psi_2 \) is surjective on \( \{ u \in \mathbb{R}^{J \times K \times K} : u(j, k_1, k_2) = 0 \text{ if } j \notin \mathcal{S}_d(k_1, k_2) \} \), all the properties of a mirror map are preserved. Moreover, the corresponding Bregman divergences \( B_{\Psi_1} \) and \( B_{\Psi_2} \) are given by

\[
B_{\Psi_1}(v_1, v_2) = \sum_i v_1(i) \log \frac{v_1(i)}{v_2(i)} - v_1(i) + v_2(i)
\]

and

\[
B_{\Psi_2}(u_1, u_2) = \sum_{j, k_1, k_2} 1(j \in \mathcal{S}_d(k_1, k_2)) \left[ u_1(j, k_1, k_2) \log \frac{u_1(j, k_1, k_2)}{u_2(j, k_1, k_2)} - u_1(j, k_1, k_2) + u_2(j, k_1, k_2) \right].
\]

(176)

(177)

Note that when \( v_1, v_2 \in \mathcal{E}_1 \) the Bregman divergence \( B_{\Psi_1}(v_1, v_2) \) reduces to relative entropy between \( v_1 \) and \( v_2 \). Similarly, when \( u_1, u_2 \in \mathcal{E}_2 \) the Bregman divergence \( B_{\Psi_2}(u_1, u_2) \) reduces to relative entropy between \( u_1 \) and \( u_2 \). It is known that the Bregman divergence (176) is convex in both arguments (17). Similarly, (177) is convex in both arguments as well.

We will use the algorithm mirror prox [18], known to be able to iteratively find the saddlepoint for convex-concave optimization problems where the gradients \( \nabla_v f(v, u) \) and \( \nabla_u f(v, u) \) are Lipschitz functions. Unfortunately, this condition does not hold in our problem, because the gradient is not necessarily finite on the boundaries of both \( \mathcal{E}_1, \mathcal{E}_2 \). Therefore, first we prove the following result to control the growth of the gradient. Then using this fact, we add an additional step to the standard mirror prox algorithm and show that it converges to the saddlepoint.

**Lemma 8:** Let \( v_0, u_0 \) be defined as

\[
v_0(j) = \frac{1}{J}, \quad \forall j \in \mathcal{X}
\]

and

\[
u_0(j, k_1, k_2) = \frac{W(k_1 | j) I(j \in \mathcal{S}_d(k_1, k_2))}{\sum_{k_2} I(j \in \mathcal{S}_d(k_1, k_2))}, \quad \forall (j, k_1, k_2) \in \mathcal{X} \times \mathcal{Y} \times \mathcal{Y}
\]

(178)

(179)

Let \( \kappa = \frac{1}{J} \), then for any \((v', u') \in \mathcal{E}_1 \times \mathcal{E}_2\)

\[
|\nabla_{v = (1 - \kappa)v' + \kappa \nu_0} f(v, (1 - \kappa)u' + \kappa u_0)|_{\infty} \leq \log \frac{TJ}{W_{\min}} + \log(K) + 1
\]

(180)

\[
|\nabla_{u = (1 - \kappa)u' + \kappa \nu_0} f((1 - \kappa)v' + \kappa v_0, u)|_{\infty} \leq \log(TJ).
\]

(181)

where \( W_{\min} = \min_{j \in \mathcal{X}, k \in \mathcal{Y}, \frac{W(k)}{W(k_j)} > 0} W(k | j) \).

**Proof:** In the following expressions \( P_X, P_{Y | X} \) correspond to \((1 - \kappa)v' + \kappa \nu_0, u\), respectively. Note that every entry of \((1 - \kappa)v' + \kappa \nu_0\) is greater than or equal to \( \frac{1}{J} \). As a result, every entry of \( q_Y\), which is output distribution corresponding to \( P_X, P_{Y | X}\), is either 0 or greater than or equal to \( \frac{W_{\min}}{TJ} \). Therefore, (180) follows by

\[
\left| \frac{\partial}{\partial P_X(j)} f(P_X, P_{Y | X}) \right| = \left| \sum_{k \in \mathcal{Y}} P_{Y | X}(k | j) \log \frac{P_{Y | X}(k | j)}{q_Y(k)} - 1 \right|
\]

(182)

\[
= \left| -H(\hat{Y} | X = j) - \log(q_Y) - 1 \right|
\]

(183)

\[
\leq \log(K) + \log \frac{TJ}{W_{\min}} + 1
\]

(184)
For (181) we have when \( j \in S_d(k_1, k_2) \) then, \( P_{\hat{Y}|X}(k_2|j) \geq P_{Y|X}(k_1, k_2|j) \geq \frac{W_{\text{min}}}{W_{JK}} \), therefore

\[
\frac{\partial}{\partial P_{Y|X}(k_1, k_2|j)} f(P_X, P_{Y|X}) \bigg|_{P_X} = \left| P_X(j) \log \frac{P_{\hat{Y}|X}(k_2|j)}{q_{\hat{Y}}(k_2)} \right| \tag{185}
\]

\[
\leq \left| \log P_{\hat{Y}|X}(k_2|j) \right| + \left| \log q_{\hat{Y}}(k_2) \right| \tag{186}
\]

\[
\leq \log \frac{T K J}{W_{\text{min}}} \log \frac{T J}{W_{\text{min}}} \tag{187}
\]

For ease of notation, let \( G = \max \{ \log(K) + \log \frac{T J}{W_{\text{min}}} \} + 1, \log \frac{T J}{W_{\text{min}}} \} \) in the rest of the paper. \( \blacksquare \)

From the choices of \( v_0, u_0 \) in (178) and (179), and (176) and (177) we get that for all \( v \in \mathcal{E}_1 \) and \( u \in \mathcal{E}_2 \)

\[
\max_{v \in \mathcal{E}_1} B_{\Psi_1}(v, v_0) \leq \log(J) \tag{188}
\]

\[
\max_{u \in \mathcal{E}_2} B_{\Psi_2}(u, u_0) \leq \log \frac{JK^2}{W_{\text{min}}} \tag{189}
\]

Moreover, (188) is true since the relative entropy between a distribution and the uniform distribution is bounded by the logarithm of the alphabet cardinality. Moreover, by definition (179) all of the entries of \( u_0 \) are either 0 or not less than \( \frac{W_{\text{min}}}{JK} \). Additionally, by definition of set \( \mathcal{E}_2 \), \( u \) is equal to zero at entries that \( u_0 \) equals to zero. Using these two facts (189) follows. Now let \( v_t \) and \( u_t \), \( t = 1, 2, \ldots, T \) be defined by the following iterative algorithm, where \( T \) is the maximum number of iterations. Let the initial values \( v_0, u_0 \) be chosen from (178) and (179), respectively. The algorithm computes \( v_{t+1} \) as follows

\[
\nabla \Psi_1(\tilde{v}_{t+1}) = \nabla \Psi_1(v_t) - \eta_t \nabla_{v=v_t} f(v, u_t) \tag{190}
\]

\[
v_{t+1}' = \arg \min_{v \in \mathcal{E}_1} B_{\Psi_1}(v, \tilde{v}_{t+1}) \tag{191}
\]

\[
v_{t+1} = (1 - \kappa)v_{t+1}' + \kappa v_0 \tag{192}
\]

where \( \eta_t \) is the stepsize at iteration \( t \). From the definition of mirror map, the range of \( \nabla \Psi_1 \) is \( R^J \), guaranteeing the existence of \( \tilde{v}_{t+1} \) in (190). Similarly, the algorithm computes \( u_{t+1} \) as follows

\[
\nabla \Psi_2(\tilde{u}_{t+1}) = \nabla \Psi_2(u_t) - \eta_t \nabla_{u=u_t} f(v_t, u) \tag{193}
\]

\[
u_{t+1}' = \arg \min_{u \in \mathcal{E}_2} B_{\Psi_2}(u, \tilde{u}_{t+1}) \tag{194}
\]

\[
u_{t+1} = (1 - \kappa)u_{t+1}' + \kappa u_0 \tag{195}
\]

Similarly, the range of \( \nabla \Psi_2 \) is \( \left\{ u \in \mathbb{R}^{JK \times K} \bigg| u(j, k_1, k_2) = 0 \text{ if } j \notin S_d(k_1, k_2) \right\} \), guaranteeing the existence of \( \tilde{u}_{t+1} \) in (193).

The following result guarantees the convergence of proposed iterative algorithm to the saddlepoint.

**Proposition 1:** Let \( \kappa = \frac{1}{T} \) and the stepsize \( \eta_t = \sqrt{T} \). Then, we have:

\[
\left| f \left( \frac{1}{T} \sum_{t=1}^T v_t, \frac{1}{T} \sum_{t=1}^T u_t \right) - \max_{v \in \mathcal{E}_1} \min_{u \in \mathcal{E}_2} f(v, u) \right| \leq \frac{1}{\sqrt{T}} \left( 4 \log \frac{JK^2}{W_{\text{min}}} + G^2 \right) \tag{196}
\]

**Proof:** We assume several properties of Bregmann divergences without proof. For further details see (17).

The first-order optimality of Bregmann divergence projections states that for any \( v_* \in \mathcal{E}_1 \) (17)

\[
\langle g_t, v_{t+1}' - v_* \rangle \leq \frac{1}{\eta_t} \left( B_\varphi(v_*, v_t) - B_\varphi(v_*, v_{t+1}') - B_\varphi(v_{t+1}', v_t) \right), \tag{197}
\]
where for ease of notation we have defined \( g_t \triangleq \nabla_{v_t} f(v, u_t) \). Now for any arbitrary \( v_0 \in E_1 \) we have

\[
\sum_{t=1}^{T} [f(v_t, u_t) - f(v_0, u_t)] \leq \sum_{t=1}^{T} (g_t, v_t - v_0) \leq \sum_{t=1}^{T} (|g_t| + (g_t, v_t - v_0)) \leq \sum_{t=1}^{T} \left( \frac{1}{2\eta_t} |v_t - v_0|^2 + \frac{\eta_t G^2}{2} \right)
\]

(198)

(199)

(200)

(201)

(202)

where (198) follows from the definition of convexity of \( f \), (201) follows from Hölder’s inequality (19) \( (g_t, v_t - v_0) \leq |v_t - v_0|_2 |g_t|_2 \) and \( |g_t|_2 \leq G \) from Lemma 8. Moreover, inequality (202) follows from (199). Now from convexity of \( B_{\psi_1} (\cdot, \cdot) \) in the second argument we have that

\[
B_{\psi_1}(v_0, v_t) = (1 - \kappa)B_{\psi_1}(v_0, v'_t) + \kappa B_{\psi_1}(v_0, v_t).
\]

(203)

Therefore plugging (203) in (202) we get:

\[
\sum_{t=1}^{T} (f(v_t, u_t) - f(v_0, u_t)) \leq \frac{1}{\eta_t} (B_{\psi_1}(v_0, v_t) + T \kappa B_{\psi_1}(v_0, v_0)) + \frac{T \eta_t G^2}{2}
\]

(204)

Now setting \( \kappa = \frac{1}{T}, \eta_t = \sqrt{T} \) and noticing \( B_{\psi_1}(v, v_0) \leq \log J \leq \log \frac{JK^2}{W_{\min}} \) we get

\[
\frac{1}{T} \left( \sum_{t=1}^{T} f(v_t, u_t) - f(v_0, u_t) \right) \leq \frac{1}{\sqrt{T}} \left( 2 \log \frac{JK^2}{W_{\min}} + \frac{1}{2} G^2 \right)
\]

(205)

The same procedure gives

\[
\frac{1}{T} \left( \sum_{t=1}^{T} f(v_t, u_t) - f(v_0, u_0) \right) \geq \frac{-1}{\sqrt{T}} \left( 2 \log \frac{JK^2}{W_{\min}} + \frac{1}{2} G^2 \right).
\]

(206)

Now we have

\[
f \left( \frac{1}{T} \sum_{t=1}^{T} v_t, u_t \right) - f \left( v_*, \frac{1}{T} \sum_{t=1}^{T} u_t \right) \leq \frac{1}{T} \left( \sum_{t=1}^{T} f(v_t, u_t) - f(v_*, u_t) \right)
\]

(207)

\[
\leq \frac{1}{\sqrt{T}} \left( \log J + G^2 \right)
\]

(208)

where (207) follows from the convex-concave nature of \( f \) and (208) follows from summing (205) and (206).

Since \( v_0 \) and \( u_0 \) were arbitrary, let \( v_* = \arg \min_{v \in E_1} f \left( v, \frac{1}{T} \sum_{t=1}^{T} u_t \right) \) and \( u_* = \arg \max_{u \in E_2} f \left( \frac{1}{T} \sum_{t=1}^{T} v_t, u \right) \) then we have:

\[
f \left( v_*, \frac{1}{T} \sum_{t=1}^{T} u_t \right) = \min_{v \in E_1} f \left( v, \frac{1}{T} \sum_{t=1}^{T} u_t \right) \leq \min_{v \in E_1} \max_{u \in E_2} f(v, u) \leq \max_{u \in E_2} f \left( \frac{1}{T} \sum_{t=1}^{T} v_t, u \right) = f \left( \frac{1}{T} \sum_{t=1}^{T} v_t, u_* \right)
\]

(209)
In addition, observe that
\[
f\left( v_t, \frac{1}{T} \sum_{t=1}^{T} u_t \right) \leq f\left( \frac{1}{T} \sum_{t=1}^{T} v_t, \frac{1}{T} \sum_{t=1}^{T} u_t \right) \leq f\left( \frac{1}{T} \sum_{t=1}^{T} v_t, u_t \right),
\]
and therefore by combining (209) and (210) we have
\[
\left| f\left( \frac{1}{T} \sum_{t=1}^{T} v_t, \frac{1}{T} \sum_{t=1}^{T} u_t \right) - \min_{v \in \mathcal{E}_1} \max_{u \in \mathcal{E}_2} f(v, u) \right| \leq f\left( \frac{1}{T} \sum_{t=1}^{T} v_t, u_t \right) - f\left( v_t, \frac{1}{T} \sum_{t=1}^{T} u_t \right)
\]
(211)
This combined with (208) finishes the proof.

It is well known that the divergence projection step (191) can be computed efficiently as [17], (20)
\[
v_{t+1}'(j) = \frac{\tilde{v}_{t+1}(j)}{\sum_{j'} \tilde{v}_{t+1}(j')},
\]
Similarly, the divergence projection on \( \mathcal{M}_{\text{max}}(d) \cap \{P_{Y|X}|P_{Y|X} = W\} \) in step (194) can be computed efficiently as
\[
u_{t+1}'(j, k_1, k_2) = \frac{W(k_1|j)\tilde{y}_{t+1}(j, k_1, k_2)|j \in \mathcal{S}(k_1, k_2))}{\sum_{k_2} \tilde{y}_{t+1}(j, k_1, k_2)|j \in \mathcal{S}(k_1, k_2))}.
\]
Then the iterative algorithm at iteration \( t+1 \) is summarized by the following:
\[
\tilde{v}_{t+1} = v_t \odot \exp\left( \frac{1}{\sqrt{T}} \nabla_{v=v_t} f(v_t, u_t) \right)
\]
\[
v_{t+1}'(j) = \frac{\tilde{v}_{t+1}(j)}{\sum_{j'} \tilde{v}_{t+1}(j')},
\]
\[
v_{t+1} = \frac{T-1}{T} v_{t+1}' + \frac{1}{T} v_0.
\]
and
\[
\tilde{u}_{t+1} = u_t \odot \exp\left( \frac{1}{\sqrt{T}} \nabla_{u=u_t} f(v_t, u_t) \right)
\]
\[
u_{t+1}'(j, k_1, k_2) = \frac{W(k_1|j)\tilde{u}_{t+1}(j, k_1, k_2)|j \in \mathcal{S}(k_1, k_2))}{\sum_{k_2} \tilde{u}_{t+1}(j, k_1, k_2)|j \in \mathcal{S}(k_1, k_2))}
\]
\[
u_{t+1} = \frac{T-1}{T} u_{t+1}' + \frac{1}{T} u_0,
\]
where \( a \odot b \) denotes componentwise product of the entries of vectors \( a, b \), (214) corresponds to the step (190) and (215) corresponds to the step (191). Now we can use the gradients computed in the previous section to run the algorithm. Figure 1 illustrates the convergence of \( \hat{R}_d(W) \) over the iterations \( t \) to the upper bound \( \bar{R}_d(W) \), using the suggested iterative algorithm for the channel and metric of Example 2. For reference, we also plot the \( C(W) \) and the LM rate \( R^\text{LM}_d(W) \). We have chosen an equiprobable maximal joint conditional distribution as initial condition and the initial step size \( \eta_{t=1} = \frac{1}{\sqrt{T}} \).

X. MULTILETTER BOUND

In this section, we study the multiletter extension of the bound (17). In particular, we show that the multiletter version cannot improve on the single-letter bound. We define the \( \ell \)-letter decoding metric \( d^{(\ell)} : X^\ell \times Y^\ell \to \mathbb{R} \) as follows
\[
d^{(\ell)}((x_1, x_2, \ldots, x_\ell), (y_1, y_2, \ldots, y_\ell)) = \sum_{i=1}^{\ell} d(x_i, y_i).
\]
(221)
Iteration number $t$

Rate (bits/channel use)

Fig. 1. Convergence of the proposed iterative algorithm to compute $\bar{R}_d(W)$ for the channel and metric from Example 2.

This decoding metric definition is consistent with the additive decoder we have defined in (5). We denote $j \in \mathcal{X}^\ell$ and $k \in \mathcal{Y}^\ell$ as the $\ell$-letter inputs and outputs, respectively. Let $W^{(\ell)}$ denote a DMC over input alphabet $\mathcal{X}^\ell$ and output alphabet $\mathcal{Y}^\ell$ with the channel rule $W^{(\ell)}((y_1, y_2, \cdots, y_\ell)|(x_1, x_2, \cdots, x_\ell)) = \prod_{i=1}^{\ell} W(y_i|x_i)$. Additionally, we define $P_X^{(\ell)}$ and $P_{YY|X}^{(\ell)}$ accordingly

$$P_X^{(\ell)}(x_1, \cdots x_\ell) = \prod_{i=1}^{\ell} P_X(x_i)$$

$$P_{YY|X}^{(\ell)}((y_1, y_2, \cdots, y_\ell), (\hat{y}_1, \hat{y}_2, \cdots, \hat{y}_\ell)|(x_1, x_2, \cdots, x_\ell)) = \prod_{i=1}^{\ell} P_{Y|X}(y_i, \hat{y}_i|x_i)$$

$X^\ell$ and $Y^\ell$, $\hat{Y}^\ell$ denote random variables defined on alphabets $\mathcal{X}^\ell$, $\mathcal{Y}^\ell$ and $\hat{\mathcal{Y}}^\ell$, respectively. Moreover, $S_d^{(\ell)}(k_1, k_2)$ is defined as

$$S_d^{(\ell)}(k_1, k_2) \triangleq \{ i \in \mathcal{X}^\ell \mid i = \arg \max_{i' \in \mathcal{X}^\ell} d^{(\ell)}(i', k_2) - d^{(\ell)}(i', k_1) \}.$$  

In the following lemma we characterize the sets $S_d^{(\ell)}(k_1, k_2)$ and relate them to $S_d(k_{1,i}, k_{2,i}), i = 1, 2, \cdots, \ell$.

**Lemma 9:** For $j \in \mathcal{X}^\ell$, $k_1 \in \mathcal{Y}^\ell$, $k_2 \in \mathcal{Y}^\ell$ we have that $j \in S_d^{(\ell)}(k_1, k_2)$ if and only if for all $1 \leq i \leq \ell$ we have

$$j_i \in S_d(k_{1,i}, k_{2,i}).$$
Proof: We have
\[
\arg \max_{j \in \mathcal{X}^\ell} d^{(\ell)}(j, k_2) - d^{(\ell)}(j, k_1) = \arg \max_{j \in \mathcal{X}^\ell} \sum_{i=1}^\ell d(j_i, k_{2,i}) - d(j_i, k_{2,i}) \tag{226}
\]
\[
= \arg \max_{(j_1, j_2, \cdots, j_\ell) \in \mathcal{X}^\ell} \sum_{i=1}^\ell d(j_i, k_{2,i}) - d(j_i, k_{2,i}) \tag{227}
\]
From (227) we get that if \((j_1, j_2, \cdots, j_\ell) \in S_d(k_1, k_2)\) then for all \(1 \leq i \leq \ell\) we should have \(j_i \in S_d(k_{1,i}, k_{2,i})\). Therefore,
\[
S_d^{(\ell)}(k_1, k_2) = S_d(k_{1,1}, k_{2,1}) \times S_d(k_{1,2}, k_{2,2}) \times \cdots \times S_d(k_{1,\ell}, k_{2,\ell}). \tag{228}
\]

For the above \(\ell\)-letter alphabets and distributions, the construction and analysis of the bound remains unchanged. Therefore, (17) remains valid for its \(\ell\)-letter extension, which can be written as
\[
\tilde{R}_d^{(\ell)}(W) \triangleq \frac{1}{\ell} \tilde{R}_{d^{(\ell)}}(W^{(\ell)}) \tag{229}
\]
\[
= \frac{1}{\ell} \max_{P_X^{\ell}, P_{Y^{(\ell)}|X^{\ell}} \in \mathcal{M}_{\ell}(d^{(\ell)})} \min_{P_{Y^{(\ell)}|X^{(\ell)}=W^{(\ell)}}} I(P_{X^{(\ell)}}, P_{Y^{(\ell)}|X^{(\ell)}}). \tag{230}
\]
We have the following result.

Proposition 2:
\[
\tilde{R}_d^{(\ell)}(W) = \tilde{R}_d(W). \tag{231}
\]

Proof: Given that \(I(P_X, P_{Y|X})\) is convex in \(P_{Y|X}\), and concave in \(P_X\), the KKT conditions are also sufficient for global optimality. Similarly, \(I(p_{X^{(\ell)}}, P_{Y^{(\ell)}|X^{(\ell)}})\) is convex in terms of \(p_{X^{(\ell)}}\) and concave in terms of \(P_{Y^{(\ell)}|X^{(\ell)}}\). Here we use the optimality conditions derived in the previous section to show that if \(P_{X^{(\ell)}}^*, P_{Y_{Y|X}^{(\ell)}}^*\) are the optimal distributions for the single letter bound then \(P_{X^{(\ell)}}^*, P_{Y_{Y|X}^{(\ell)}}^*\) are optimal distributions for the multiletter version. If we find a feasible pair \(P_{Y^{(\ell)}|X^{(\ell)}}, P_{X^{(\ell)}}\) such that when fixing \(P_{Y^{(\ell)}|X^{(\ell)}}, \) the input distribution \(P_{X^{(\ell)}}\) is a maximizer of \(f(\cdot, P_{Y^{(\ell)}|X^{(\ell)}})\), and when fixing \(P_{X^{(\ell)}}\), the joint conditional distribution \(P_{Y^{(\ell)}|X^{(\ell)}}\) is a minimizer of \(f(p_{X^{(\ell)}}, \cdot)\), then the pair \((P_{X^{(\ell)}}, P_{Y^{(\ell)}|X^{(\ell)}})\) is a saddlepoint.

We need to show that if \(P_{X^{(\ell)}}, P_{Y^{(\ell)}|X}^*\) is a saddlepoint for the single-letter case, then \(P_{X^{(\ell)}}, P_{Y^{(\ell)}|X}^*\) is a saddlepoint for the multiletter bound. Therefore, we must show that if \(P_{X^{(\ell)}}, P_{Y^{(\ell)}|X}^*\) is a saddlepoint for the multiletter bound, then it must happen that \(P_{X^{(\ell)}} = P_{X^{(\ell)}}, P_{Y^{(\ell)}|X^{(\ell)}} = P_{Y^{(\ell)}|X^{(\ell)}}^*\). Based on the aforementioned argument, it is sufficient to show that \(P_{Y^{(\ell)}|X}^*\) is a minimizer of (230) by fixing \(P_{X^{(\ell)}}^*\). This is because it is known that \(\frac{1}{\ell} C(P_{Y^{(\ell)}|X}^*) = C(P_{X^{(\ell)}}, i.e. the product distribution \(P_{X^{(\ell)}}\) achieves \(C(P_{Y^{(\ell)}|X}^*)\).

In the following lemma we prove that by fixing \(P_{X^{(\ell)}}\), then \(P_{Y^{(\ell)}|X}\) satisfies the KKT conditions and hence, it is a minimizer of (230). Before stating the result we recall that the multiletter counterparts of the single-letter KKT conditions given in (163) and (164) hold. Moreover, as in the single-letter case, the multiletter KKT conditions are sufficient for global optimality, because the function \(f(P_{X^{(\ell)}}, \cdot)\) is concave. Using Lemma 10 below completes the proof.
Lemma 10: Let \( P^*_{XX',Y|X} \) be a saddlepoint for optimization problem \( 17 \). Set \( P_X = P^{* (t)}_X \). Then, the joint conditional distribution \( P^{* (t)}_{Y|X} \) is a minimizer of
\[
\min_{P_{Y|X} \in \mathcal{M}_{\text{max}}(d^{(t)})} I\left(P^{* (t)}_X, P^{* (t)}_{Y|X}^t\right).
\]

\( \text{(232)} \)

Proof: We should show that by setting \( P_X = P^{* (t)}_X \), the multiletter versions of the KKT conditions \( 163 \) and \( 164 \) hold for \( P^{* (t)}_{Y|X} \). Generalizing the conditions of \( 163 \) and \( 164 \) to the multi-letter case, and setting \( P_{Y|X} = P^{* (t)}_{Y|X} \), we should show that for all \( j, k_1 \in X^* \times Y^* \) there exist multipliers \( \lambda_{j,k_1} \) such that the conditions below are fulfilled. If we show this, then the Lemma is proved because these are precisely the conditions for the minimizer of \( 232 \).

1) When \( P^{* (t)}_{Y|X}(k_1, k_2|j) > 0 \) we must have,
\[
\frac{\partial}{\partial P_{Y|X}(k_1, k_2|j)} f(P^{* (t)}_X, P^{* (t)}_{Y|X}) \bigg|_{P_{Y|X} = P^{* (t)}_{Y|X}} = \lambda_{j,k_1}.
\]
\( \text{(233)} \)

2) When \( P^{* (t)}_{Y|X}(k_1, k_2|j) = 0 \) and \( j \in S_d^{(t)}(k_1, k_2) \) we must have that,
\[
\frac{\partial}{\partial P_{Y|X}(k_1, k_2|j)} f(P^{* (t)}_X, P^{* (t)}_{Y|X}) \bigg|_{P_{Y|X} = P^{* (t)}_{Y|X}} \geq \lambda_{j,k_1}.
\]
\( \text{(234)} \)

Similarly to \( 160 \), the derivative in \( 233 \) and \( 234 \) is given by,
\[
\frac{\partial}{\partial P_{Y|X}(k_1, k_2|j)} f(P^{* (t)}_X, P^{* (t)}_{Y|X}) \bigg|_{P_{Y|X} = P^{* (t)}_{Y|X}} = P^{* (t)}_X(j) \log \frac{P^{* (t)}_{Y|X}(k_1|j)}{q^{* (t)}_Y(k_1)}
\]
\( \text{(235)} \)

which, by using that in \( P_{Y|X} = P^{* (t)}_{Y|X} \) and \( q^{* (t)}_Y \) are product distributions, gives,
\[
P^{* (t)}_X(j) \log \frac{P^{* (t)}_{Y|X}(k_1|j)}{q^{* (t)}_Y(k_1)} = P^*_X(j_1)P^*_X(j_2) \cdots P^*_X(j_\ell)
\]
\[
\left( \log \frac{P^{* (t)}_{Y|X}(k_2,1|j_1)}{q^{* (t)}_Y(k_2,1)} + \log \frac{P^{* (t)}_{Y|X}(k_2,2|j_2)}{q^{* (t)}_Y(k_2,2)} + \cdots + \log \frac{P^{* (t)}_{Y|X}(k_2,\ell|j_\ell)}{q^{* (t)}_Y(k_2,\ell)} \right)
\]
\( \text{(236)} \)

In order to show that there exist some coefficients \( \lambda_{j,k_1} \) satisfying both \( 233 \) and \( 234 \), we make a particular choice and show that this choice satisfies both \( 233 \) and \( 234 \). To this end, define
\[
\lambda_{j,k_1} = \begin{cases} 
0 & \text{if } P^*_X(j_1)P^*_X(j_2) \cdots P^*_X(j_\ell) = 0 \\
P^*_X(j_1)P^*_X(j_2) \cdots P^*_X(j_\ell) \left( \frac{\lambda_{j_1,k_1}}{P^*_X(j_1)} + \frac{\lambda_{j_2,k_1}}{P^*_X(j_2)} + \cdots + \frac{\lambda_{j_\ell,k_1}}{P^*_X(j_\ell)} \right) & \text{if } P^*_X(j_1)P^*_X(j_2) \cdots P^*_X(j_\ell) \neq 0 
\end{cases}
\]
\( \text{(237)} \)

where \( \lambda_{j_1,k_1} \) is the single-letter Lagrange multiplier corresponding to \( j_1 \) and \( k_1,j_1 \).

Now, excluding the cases where \( P^*_X(j_1)P^*_X(j_2) \cdots P^*_X(j_\ell) = 0 \) where from \( 236 \), \( 233 \) and \( 234 \) the KKT conditions clearly hold, we have two cases

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1) When $P_{Y|X}^*(j, k_1, k_2) > 0$, then for all $1 \leq i \leq \ell$ we must have $P_{Y|X}^*(k_{1,i}, k_{2,i}|j_i) > 0$ and therefore, (163) is valid. We have to verify that this implies that (233) is also valid. Thus,
\[
\frac{\partial}{\partial P_{Y^*|X^*}(k_1, k_2)} f(P_{X^*}, P_{Y^*|X^*}) \bigg| _{P_{Y^*|X^*} = P_{Y|X}^*} = P_{X^*}^*(j_1)P_{X^*}^*(j_2) \cdots P_{X^*}^*(j_i) \left( \log \frac{P_{X|X}^*(k_{1,i}|j_i)}{q_{Y_i}^*} + \log \frac{P_{Y|X}^*(k_{2,i}|j_i)}{q_{Y_i}^*} + \cdots + \log \frac{P_{Y|X}^*(k_{2,i}|j_i)}{q_{Y_i}^*} \right) 
\]
\[
= P_{X^*}^*(j_1)P_{X^*}^*(j_2) \cdots P_{X^*}^*(j_i) \left( \frac{\lambda_{j_1,k_{1,i}}}{P_{X^*}^*(j_i)} + \frac{\lambda_{j_2,k_{2,i}}}{P_{X^*}^*(j_i)} + \cdots + \frac{\lambda_{j_{\ell},k_{\ell,i}}}{P_{X^*}^*(j_i)} \right) 
\]
\[
= \lambda_{j,k_i}
\]  
where (240) holds from the single-letter optimality in (163).

2) When $P_{Y|X}^*(k_1, k_2|j) = 0$ and $j \in S_{d^*}(k_1, k_2)$, as a result of Lemma, we have that $S_{d^*}(k_1, k_2)$ is a product set, i.e. for all $1 \leq i \leq \ell$, 
\[
ji \in S_{d}(k_{1,i}, k_{2,i}).
\]
Moreover, either $P_{Y|X}^*(k_{1,i}, k_{2,i}|j_i) > 0$ where (163) is satisfied or $P_{Y|X}^*(k_{1,i}, k_{2,i}|j_i) = 0$ where (164) is satisfied. Now, with these assumptions, we should verify that (234) is valid. We have,
\[
\frac{\partial}{\partial P_{Y^*|X^*}(k_1, k_2)} f(P_{X^*}, P_{Y^*|X^*}) \bigg| _{P_{Y^*|X^*} = P_{Y|X}^*} = P_{X^*}^*(j_1)P_{X^*}^*(j_2) \cdots P_{X^*}^*(j_i) \left( \log \frac{P_{X|X}^*(k_{1,i}|j_i)}{q_{Y_i}^*} + \log \frac{P_{Y|X}^*(k_{2,i}|j_i)}{q_{Y_i}^*} + \cdots + \log \frac{P_{Y|X}^*(k_{2,i}|j_i)}{q_{Y_i}^*} \right) 
\]
\[
\geq P_{X^*}^*(j_1)P_{X^*}^*(j_2) \cdots P_{X^*}^*(j_i) \left( \frac{\lambda_{j_1,k_{1,i}}}{P_{X^*}^*(j_i)} + \frac{\lambda_{j_2,k_{2,i}}}{P_{X^*}^*(j_i)} + \cdots + \frac{\lambda_{j_{\ell},k_{\ell,i}}}{P_{X^*}^*(j_i)} \right) 
\]
\[
= \lambda_{j,k_i}
\]  
where (245) is true because of the single-letter optimality in (163) and (164).

**APPENDIX A**

In this Appendix we provide the proof of Theorem 2. Without loss of generality, we assume that the sequence $\{d(1,k) - d(2,k)\}_{k=1}^K$ is non-decreasing, i.e. for $k_1 \leq k_2$,
\[
d(1,k_1) - d(2,k_1) \leq d(1,k_2) - d(2,k_2).
\]  
(247)

We can assume this, since it is always possible to relabel the output alphabet such that this property is fulfilled. This assumption simplifies the evaluation of the sets $S_{\{\cdot, \cdot\}}$. For $k_1 = k_2$ we have $S(k_1, k_2) = \{1, 2\}$. Moreover, when $k_1 < k_2$ from (247) and Definition 1, we have that $1 \in S(k_1, k_2)$ and $2 \in S(k_2, k_1)$.

We prove a slightly stronger result. In particular, we prove that the condition $C_d(W) = C(W)$ implies that sequences
\[
\left\{ P_{X^*}^*(k) \log \frac{W(k|1)}{q_Y^*(k)} \right\}_{k=1}^K, \left\{ -P_{X^*}^*(k) \log \frac{W(k|2)}{q_Y^*(k)} \right\}_{k=1}^K
\]  
(248)
both should have the same order as the decoding metric difference sequence \(\{d(1, k) - d(2, k)\}_{k=1}^{K}\), where the notation \(P_X^*\) refers to the capacity-achieving distribution of \(W\); \(q^*\) denotes the corresponding output distribution.

Now assume that \(C_d(W) = C(W)\). Therefore, \(P_X^*, P_{Y|X} = P_{Y|X}\) must be a saddle point of (146). As a result, the KKT conditions in (163) and (164) must hold. Observe that

\[
P_{Y|X}(k_1, k_2) = \begin{cases} \frac{W(k_1|j)}{q_Y(k_1)} & k_1 = k_2 \\ 0 & k_1 \neq k_2. \end{cases}
\]

Therefore, combining the KKT conditions in (163) and (164) we obtain,

1) If \(k_1 = k_2\), for both \(j = 1, 2\) we have

\[
P_X^*(j) \log \frac{W(k_1|j)}{q_Y(k_1)} = \lambda_{j,k_1}
\]

2) If \(k_1 < k_2\) we know \(1 \in S(k_1, k_2)\) and \(2 \in S(k_2, k_1)\)

\[
P_X^*(1) \log \frac{W(k_2|1)}{q_Y(k_2)} \geq \lambda_{k_1}
\]

\[
P_X^*(2) \log \frac{W(k_1|2)}{q_Y(k_1)} \geq \lambda_{k_2}
\]

Therefore we get if \(k_1 < k_2\)

\[
P_X^*(1) \log \frac{W(k_2|1)}{q_Y(k_2)} \geq \lambda_{k_1} = P_X^*(1) \log \frac{W(k_1|1)}{q_Y(k_1)}
\]

\[
P_X^*(2) \log \frac{W(k_1|2)}{q_Y(k_1)} \geq \lambda_{k_2} = P_X^*(2) \log \frac{W(k_2|2)}{q_Y(k_2)}.
\]

Therefore we get that \(\left\{P_X^*(1) \log \frac{W(k_1|1)}{q_Y(k_1)}\right\}_{k=1}^{K}\) and \(-\left\{P_X^*(2) \log \frac{W(k_2|2)}{q_Y(k_1)}\right\}_{k=1}^{K}\) are both non-decreasing sequences and so is any linear combination of them with positive coefficients. Therefore, since

\[
\log W(k|1) - \log W(k|2) = \frac{1}{P_X^*(1)} \left( P_X^*(1) \log \frac{W(k|1)}{q_Y(k)} \right) - \frac{1}{P_X^*(2)} \left( P_X^*(2) \log \frac{W(k|2)}{q_Y(k)} \right)
\]

we conclude that the sequence \(\left\{\log W(k|1) - \log W(k|2)\right\}_{k=1}^{K}\) is a non-decreasing sequence.

**Appendix B**

This section is about the choice of \(\epsilon\) in proof of the main theorem in Section VII. Let \(f : \mathcal{A} \to \mathbb{R}\) be a continuous function and \(\mathcal{A}\) be a compact set. Then this function is uniformly continuous. Now we can use this fact on entropy function \(H : \Delta^J \to \mathbb{R}\) where \(\Delta^J = \{x \in \mathbb{R}^J | x_i \geq 0, i = 1, 2, \cdots, J\}\) is the \(J\)-dimensional simplex. Therefore, for \(\hat{q}\) there is an \(\epsilon\) that if \(p_1, p_2 \in \Delta^J\) are \(\epsilon\)-close i.e. \(|p_1 - p_2|_\infty \leq \epsilon\) we have

\[
|H(p_1) - H(p_2)| \leq \frac{\sigma}{4}
\]

Now let \(V\) be matrix of a conditional distribution and \(V_1, V_2, \cdots, V_J\) be rows of \(V\). Then for any type \(p\) and any other conditional distribution matrix \(\hat{V}\) with rows \(\hat{V}_1, \hat{V}_2, \cdots, \hat{V}_J\) we have

\[
|p^T V - \hat{p}^T \hat{V}|_\infty \leq |V - \hat{V}|_\infty
\]

\[
|V_1 - \hat{V}_1|_\infty \leq |V - \hat{V}|_\infty.
\]
As a result, if $|V - \hat{V}|_\infty \leq \epsilon$ we get

$$|H(p^T V) - H(p^T \hat{V})|_\infty \leq \frac{\sigma}{4}$$  \hspace{1cm} (259)

$$|H(V) - H(\hat{V})| \leq \frac{\sigma}{4}.  \hspace{1cm} (260)$$

As for $H(V|p)$ we have

$$H(V|p) = \sum_{j=1}^{J} p(j) H(V_j),$$  \hspace{1cm} (261)

and thus,

$$|H(V|p) - H(\hat{V}|p)|_\infty \leq \sum_{j=1}^{J} p(i)|H(V_j) - H(\hat{V}_j)|_\infty$$  \hspace{1cm} (262)

$$\leq \frac{\sigma}{4}.  \hspace{1cm} (263)$$

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