Robustness of Pisot-regular sequences

Émilie Charlier
Department of Mathematics
University of Liège
Allée de la Découverte 12
4000 Liège, Belgium

Célia Cisternino∗
Department of Mathematics
University of Liège
Allée de la Découverte 12
4000 Liège, Belgium

Manon Stipulanti
Department of Mathematics
Hofstra University
Roosevelt Hall
Hempstead, NY 11549, USA

June 22, 2020

Abstract

We consider numeration systems based on a d-tuple $U = (U_1, \ldots, U_d)$ of sequences of integers and we define $(U, K)$-regular sequences through $K$-recognizable formal series, where $K$ is any semiring. We show that, for any d-tuple $U$ of Pisot numeration systems and any commutative semiring $K$, this definition does not depend on the greediness of the $U$-representations of integers. The proof is constructive and is based on the fact that the normalization is realizable by a $2d$-tape finite automaton. In particular, we use an ad hoc operation mixing a $2d$-tape automaton and a $K$-automaton in order to obtain a new $K$-automaton.

2010 Mathematics Subject Classification: 68Q70, 20M35, 11A67, 11B85.
Keywords: Regular sequence, recognizable formal series, weighted automaton, semiring, numeration system, Pisot number, normalizer

1 Introduction

A $b$-automatic sequence is a sequence $f: \mathbb{N} \to \Delta$, where $\Delta$ is a finite alphabet, that is generated by a deterministic finite automaton with output (DFAO) as follows: the $n$th term $f(n)$ of the sequence is output by the DFAO when the input is the $b$-expansion of $n$. This notion originally appeared in the work of Büchi [7], but Cobham was the first to systematically study $b$-automatic sequences, or equivalently, $b$-recognizable sets of
than 1 is an algebraic integer greater than 1 systems in a natural way. Such generalizations were considered for example in \([1, 8, 9, 14]\), \((\mathbb{U}^i, r, \mathbb{M})\) for \(i = 0, \ldots, M - 1\) such that, for all words \(a_i \cdots a_0\) over the alphabet \(\{0, \ldots, b - 1\}\), the matrix product \(\lambda M_{a_i} \cdots M_{a_0} \gamma f\) equals \(f(n)\) if \(a_i \cdots a_0\) is the \(b\)-expansion of \(n\) and equals 0 otherwise. The latter condition translates in algebraic terms by saying that the formal series (with noncommutative variables) \(\sum_{w \in \text{rep}_b(\mathbb{N})} f(\text{val}_b(w))w\) is \(\mathbb{K}\)-recognizable. In the previous series, the words \(w\) run over \(\text{rep}_b(\mathbb{N})\), which is the language of the \(b\)-expansions of non-negative integers, that is, the language made of the words over the alphabet \(\{0, \ldots, b - 1\}\) not starting with the letter 0. Note that we do not impose any condition on the semiring \(\mathbb{K}\) in our definition. Regular sequences with no condition on the underlying semiring were studied by Berstel and Reutenauer \([4]\) and used in \([8, 9]\). In particular, it can be proven that a sequence with values in a semiring \(\mathbb{K}\) is \(b\)-automatic if and only if it is \((b, \mathbb{K})\)-regular and takes only finitely many values.

The notion of \((b, \mathbb{K})\)-regular sequences can be extended to more general numeration systems in a natural way. Such generalizations were considered for example in \([1, 8, 9, 13, 10]\). In this paper, we focus on the class of Pisot numeration systems. A Pisot number is an algebraic integer greater than 1 all of whose Galois conjugates have modulus less than 1. In these systems, natural numbers are expanded via an increasing base sequence \(U = (U(i))_{i \in \mathbb{N}}\) with bounded quotients \(U(i+1)/U(i)\) and that satisfies a linear recurrence relation whose characteristic polynomial is the minimal polynomial of a Pisot number. We then talk about \((U, \mathbb{K})\)-regular sequences (precise definitions will be given in Section 2). We refer to \([12]\) for a survey on Pisot numeration systems.

In the case of integer bases, it is not hard to see that a sequence \(f : \mathbb{N} \rightarrow \mathbb{K}\) is \((b, \mathbb{K})\)-regular if and only if the series \(\sum_{w \in \{0, \ldots, b - 1\}^*} f(\text{val}_b(w))w\) is \(\mathbb{K}\)-recognizable \([4]\). Proving this equivalence consists in dealing with non-greedy representations of integers. For integer base numeration systems, this comes down to adding or removing leading zeroes. However, when considering more general numeration systems, there exist non-greedy representations of another kind. For example, the words 11 and 100 both represent the integer 3 in the Zeckendorf numeration system, which is based on the sequence \(1, 2, 3, 5, 8, 13, \ldots\) of Fibonacci numbers \([19]\). Our aim is to provide characterizations of regular sequences in the extended framework of Pisot numeration system that only depends on the value function. In other words, we want to characterize the family of \((U, \mathbb{K})\)-regular sequences in a way that does not depend on the choice of an algorithm to represent numbers. Moreover, we will carry our work in a multidimensional setting, meaning that we will define and study multidimensional \((U, \mathbb{K})\)-regular sequences \(f : \mathbb{N}^d \rightarrow \mathbb{K}\) where \(U = (U_1, \ldots, U_d)\) is a \(d\)-tuple of Pisot numeration systems.

Our method relies on the property that normalization in Pisot numeration systems is realizable by a computable finite automaton \([6, 13]\). The existence of such a normalizer will be crucial in the proof of our main result. Along the way, we will show the existence of normalizers extended both to negative values and to the multidimensional setting. Let us also stress that all our techniques are effective, meaning that linear representations are computable for all formal series that will be proven to be \(\mathbb{K}\)-recognizable.
The paper is organized as follows. In Section 2, we provide the necessary background. First, we define $K$-automata and formal series. We then introduce the notions of $K$-recognizable series and numeration systems, in particular Pisot numeration systems. In Section 3, we define multidimensional $(U, K)$-regular sequences and we state our main result, which is Theorem 8. It provides characterizations of multidimensional $(U, K)$-regular sequences that are independent of the choice of admissible $U$-representations of integers. Section 4 is concerned with normalization in Pisot numeration systems. We first recall the existence of a normalizer associated with any Pisot numeration system in dimension 1, that is, a 2-tape automaton accepting words of the form $(u^n)$ where $u$ has a non-negative value $n$ and $v$ is the greedy representation of $n$. We then extend the normalizer to words with negative values and to the multidimensional setting. Finally, in Section 5, we prove Theorem 8. In the process, we define and use a particular operation mixing 2$d$-tape automata and $K$-automata. In Section 6, we end by making the linear representations involved in the proof of Theorem 8 explicit. Throughout the paper, we use a running example in order to illustrate the definitions as well as the construction in the proof of Theorem 8.

2 Preliminaries

We make use of common notions in formal language theory, such as alphabet, letter, word, length of a word, language and usual definitions from automata theory [15]. In particular, we let $\varepsilon$ denote the empty word.

We take the convention that elements of $\mathbb{Z}^d$ are written as vertical vectors. Moreover, a word

$$
\begin{pmatrix}
ap_{1,1} \\
ap_{2,1} \\
\vdots \\
ap_{d,1} 
\end{pmatrix} \begin{pmatrix}
ap_{1,2} \\
ap_{2,2} \\
\vdots \\
ap_{d,2} 
\end{pmatrix} \cdots \begin{pmatrix}
ap_{1,\ell} \\
ap_{2,\ell} \\
\vdots \\
ap_{d,\ell} 
\end{pmatrix}
$$

in $(\mathbb{Z}^d)^*$ will be written as

$$
\begin{pmatrix}
\prod_{i=1}^{\ell} n_{1,i} \\
\prod_{i=1}^{\ell} n_{2,i} \\
\vdots \\
\prod_{i=1}^{\ell} n_{d,i}
\end{pmatrix}.
$$

In what follows, it is important to observe that if a word

$$
\begin{pmatrix}
\prod_{i=1}^{\ell} u_{1,i} \\
\prod_{i=1}^{\ell} u_{2,i} \\
\vdots \\
\prod_{i=1}^{\ell} u_{d,i}
\end{pmatrix}
$$

belongs to $(\mathbb{Z}^d)^*$ then necessarily $w_1, \ldots, w_d$ are words of equal lengths.

We use bold letters to designate $d$-dimensional elements. Most of the time, $n$ is a vector in $\mathbb{Z}^d$, $a, b$ are letters in $\mathbb{Z}^d$ and $u, v, w$ are words in $(\mathbb{Z}^d)^*$. Moreover, $0$ denotes the zero letter of dimension $d$ and $\varepsilon$ denotes the empty word of dimension $d$. Throughout the text, the context will clearly indicate if an element of $\mathbb{Z}^d$ should be seen as a vector of integers or as a letter.

In order to avoid any confusion between the notions of length and absolute value, we use the notation $|| \cdot ||$ in order to designate the component-wise absolute value of vectors of integers given by

$$
\begin{pmatrix}
\prod_{i=1}^{\ell} |n_{1,i}| \\
\prod_{i=1}^{\ell} |n_{2,i}| \\
\vdots \\
\prod_{i=1}^{\ell} |n_{d,i}|
\end{pmatrix}.
$$
We extend this notation to words over \( \mathbb{Z}^d \) by setting \( ||\varepsilon|| = \varepsilon \) and for all \( u, v \in (\mathbb{Z}^d)^* \), \( ||uv|| = ||u|| ||v|| \). The length of a word \( w \) is, as usual, written \( |w| \).

From now on, let \( \mathbb{K} \) be a fixed semiring and let \( d \) be a dimension, that is, a positive integer.

### 2.1 Weighted automata

A weighted automaton \( A = (Q, I, T, A, E) \) with weights in \( \mathbb{K} \), or simply a \( \mathbb{K} \)-automaton, is composed of a finite set \( Q \) of states, a finite alphabet \( A \) and of three mappings \( I: Q \to \mathbb{K} \), \( T: Q \to \mathbb{K} \) and \( E: Q \times A \times Q \to \mathbb{K} \). A transition \( (p, a, q) \in Q \times A \times Q \) such that \( E(p, a, q) \neq 0 \) is an edge, the letter \( a \) is its label and \( E(p, a, q) \) is its weight. With each \( q \in Q \), we associate an initial weight \( I(q) \) and a final weight \( T(q) \). We call a state \( q \in Q \) initial (resp. final) if \( I(q) \neq 0 \) (resp. \( T(q) \neq 0 \)). A path in \( A \) is a sequence \( c = (q_0, a_1, q_1)(q_1, a_2, q_2) \cdots (q_{n-1}, a_n, q_n) \) of consecutive edges. For each path \( c \), we let \( i_c \) and \( t_c \) denote the first and last states \( q_0 \) and \( q_n \) of \( c \) respectively. The weight of the path \( c \) is the product \( E(c) = E(q_0, a_1, q_1)E(q_1, a_2, q_2) \cdots E(q_{n-1}, a_n, q_n) \) of the weights of its edges and its label is the word \( a_1a_2 \cdots a_n \). For each \( w \in A^* \), we let \( C_A(w) \) denote the set of paths in \( A \) with label \( w \). The weight of a word \( w \in A^* \) in \( A \) is the quantity \( \sum_{c \in C_A(w)} I(i_c)E(c)T(t_c) \). For each path \( c \in C_A(w) \), we refer to \( I(i_c)E(c)T(t_c) \) as its contribution to the weight of \( w \).

A \( \mathbb{K} \)-automaton is represented by a graph. Each state is a vertex and each edge carries an expression of the form \( a|k \), where \( a \) is its label and \( k \) is its weight. Each initial (resp. final) state \( q \) is distinguished by an incoming (resp. outgoing) arrow which carries the weight \( I(q) \) (resp. \( T(q) \)).

**Example 1.** We start a running example by considering the \( \mathbb{N} \)-automaton over the alphabet \( \{0, 1, 2\} \) depicted in Figure 1.

![Figure 1: An N-automaton over the alphabet \( \{0, 1, 2\} \).](image)

### 2.2 Formal series

A formal series over a finite alphabet \( A \) is a function \( S: A^* \to \mathbb{K} \). The image under \( S \) of a word \( w \) is denoted by \( (S, w) \) and is called the coefficient of \( w \) in \( S \). We use the notation \( S = \sum_{w \in A^*} (S, w) \). In what follows, we simply talk about series instead of formal series.
The Hadamard product of two series $S, T : A^* \to K$ is the series $S \odot T$ defined by $\sum_{w \in A^*} (S(w), T(w))$ over $K$. If $L$ is a language over $A$, then its characteristic series is the series $L = \sum_{w \in L} w$. In particular, for $S : A^* \to K$ and $L \subseteq A^*$, we have $S \odot L = \sum_{w \in L} (S(w), T(w))$.

A series $S : A^* \to K$ is $K$-recognizable if there exist an integer $r \in \mathbb{N}_{\geq 1}$, a morphism of monoids $\mu : A^* \to K^{r \times r}$ (with respect to concatenation of words and multiplication of matrices) and two matrices $\lambda \in K^{1 \times r}$ and $\gamma \in K^{r \times 1}$ such that, for all words $w \in A^*$, $(S(w), T(w)) = \lambda \mu(w) \gamma$. In this case, the triple $(\lambda, \mu, \gamma)$ is called a linear representation of $S$.

Formal series and $K$-automata can be linked. Given a $K$-automaton $A = (Q, I, T, A, E)$, the series recognized by $A$ is the series $\sum_{w \in A^*} (\sum_{c \in C_A(w)} I(i_c)E(c)T(i_c)) w$ whose coefficients are the weights of the words in $A$.

**Proposition 2.** A series is recognized by a $K$-automaton if and only if it is $K$-recognizable.

A constructive proof can be found in [1] Chapter 1, Proposition 6.1. Roughly, the linear representation encodes the weights in the $K$-automaton.

**Example 3.** For all $a \in \{0, 1, 2\}$ and all $w \in \{0, 1, 2\}^*$, let $|w|_a$ be the number of occurrences of the letter $a$ in the word $w$. Let $g : \{0, 1, 2\}^* \to \mathbb{N}$ be the function mapping a word $w$ to the greatest $n$ such that $1^n$ is a prefix of $w$. Consider the regular language $L$ of the words over $\{0, 1, 2\}$ starting with $1$ or $2$ and avoiding factors in $21^*2$, and consider the series $S$ defined by

$$(S(w), T(w)) = \begin{cases} (3|w|_12^{|w|_2} + 3|w|_1-|g(w)|)2^{|w|_2} & \text{if } w \in L \text{ and } g(w) \neq 0 \\ 3^{|w|_1}2^{|w|_2} & \text{if } w \in L \text{ and } g(w) = 0 \\ 0 & \text{if } w \notin L \end{cases}$$

for all $w \in \{0, 1, 2\}^*$. For example, the coefficients of $1121, 2101$ and $2112$ in $S$ are $3^3 \cdot 2 + 3 \cdot 2 = 60, 3^2 \cdot 2 = 18$ and $0$ respectively. The series $S$ is $\mathbb{N}$-recognizable as a linear representation is given by $(\lambda, \mu, \gamma)$ where

$$\lambda = (10000), \mu(0) = (00000, 0100, 0100), \mu(1) = (0301, 0300, 0001), \mu(2) = (0002, 0002, 0002), \gamma = \left(\begin{array}{c} 1 \\ 1 \end{array}\right).$$

The corresponding $\mathbb{N}$-automaton, with respect to the order of the states $W, X, Y, Z$, is that of Figure [1].

### 2.3 Numeration systems

A numeration system is given by an increasing sequence $U = (U(i))_{i \in \mathbb{N}}$ of integers such that $U(0) = 1$ and $\sup_{i \in \mathbb{N}} \frac{U(i+1)}{U(i)} < +\infty$. The value function $\text{val}_U : \mathbb{Z}^* \to \mathbb{Z}$ maps any word $w = a_1 \cdots a_0 \in \mathbb{Z}^*$ to $\text{val}_U(w) = \sum_{i=0}^\ell a_i U(i)$. A $U$-representation of an integer $n$ is a word $w \in \mathbb{Z}^*$ such that $n = \text{val}_U(w)$. The greedy $U$-representation $\text{rep}_U(n)$ of a positive integer $n$ is the unique $U$-representation $a_{\ell} \cdots a_0$ of $n$ such that for all $j \in \{0, \ldots, \ell\}$, $a_j \geq 0$ and $\sum_{i=0}^j a_i U(i) < U(j+1)$, and moreover $a_\ell \neq 0$. It implies that the digits of greedy $U$-representations belong to the alphabet $A_U = \{0, \ldots, \sup_{i \in \mathbb{N}} \frac{U(i+1)}{U(i)} \} - 1$. We set $\text{rep}_U(0) = \varepsilon$. The numeration language is the set $\text{rep}_U(\mathbb{N}) = \{\text{rep}_U(n) : n \in \mathbb{N}\}$.

A numeration system is linear if it satisfies a linear recurrence relation over $\mathbb{Z}$. A Pisot number is an algebraic integer greater than 1 whose Galois conjugates all have modulus less than 1. A Pisot numeration system is a linear numeration system whose characteristic
polynomial is the minimal polynomial of a Pisot number. For such a numeration system, \( \text{rep}_U(N) \) is a regular language \([3]\). Note that Pisot numeration systems associated with the same Pisot number may only differ by their initial conditions.

**Example 4.** Consider the Pisot number \( \frac{3 + \sqrt{5}}{2} \), which is the square of the golden ratio. Its minimal polynomial is \( X^2 - 3X + 1 \). Let \( U \) be the Pisot numeration system defined by \( U(0) = 1, U(1) = 3 \) and for \( i \geq 2, U(i) = 3U(i - 1) - U(i - 2) \). This system is the Bertrand numeration system associated with \( \frac{3 + \sqrt{5}}{2} \); see \([5, 6]\). We have \( A_U = \{0, 1, 2\} \). A DFA accepting the language \( 0^*\text{rep}_U(N) \) is depicted in Figure 2. Observe that the language \( L \) considered in Example 3 is precisely the numeration language \( \text{rep}_U(N) \).

![Figure 2: An automaton accepting the language 0*rep_U(N) where U is the Bertrand numeration system associated with the Pisot number 3+√5/2.](image)

In order to be able to represent \( d \)-dimensional vectors of integers by words, we introduce the following usual convention. For a \( d \)-tuple of words
\[
\left( \begin{array}{c}
  w_1 \\
  \vdots \\
  w_d
\end{array} \right)
\]
in \( (\mathbb{Z}^*)^d \), we set
\[
\left( \begin{array}{c}
  w_1 \\
  \vdots \\
  w_d
\end{array} \right) = \left( \begin{array}{c}
  0^{\ell - |w_1|}w_1 \\
  \vdots \\
  0^{\ell - |w_d|}w_d
\end{array} \right)
\]
where \( \ell \) is the maximum of the lengths of the words \( w_1, \ldots, w_d \). Thus, we have padded the shortest words with leading zeroes in order to obtain \( d \) words of the same length. The obtained \( d \)-tuple can now be seen as a word over the \( d \)-dimensional alphabet \( \mathbb{Z}^d \), that is, an element of \( (\mathbb{Z}^d)^* \). For a \( d \)-tuple of numeration systems \( U = (U_1, \ldots, U_d) \), we set \( A_U = A_{U_1} \times \cdots \times A_{U_d} \). Accordingly, we extend the definition of the maps \( \text{rep}_U \) and \( \text{val}_U \) as follows:
\[
\text{rep}_U: \mathbb{N}^d \to A_U, \quad \left( \begin{array}{c}
  n_1 \\
  \vdots \\
  n_d
\end{array} \right) \mapsto \left( \begin{array}{c}
  \text{rep}_{U_1}(n_1) \\
  \vdots \\
  \text{rep}_{U_d}(n_d)
\end{array} \right)
\]
and
\[
\text{val}_U: (\mathbb{Z}^*)^d \to \mathbb{Z}^d, \quad \left( \begin{array}{c}
  w_1 \\
  \vdots \\
  w_d
\end{array} \right) \mapsto \left( \begin{array}{c}
  \text{val}_{U_1}(w_1) \\
  \vdots \\
  \text{val}_{U_d}(w_d)
\end{array} \right).
\]
For an easier writing, for all \( a \in \mathbb{Z} \) (considered as a letter), we write \( \overline{a} \) instead of \( -a \), and we extend this notation to words over \( \mathbb{Z} \) by setting \( \overline{uw} = \overline{w} \overline{u} \) for \( u, v \in \mathbb{Z}^* \).

**Example 5.** In dimension 2, we have \( \text{rep}_{(U,U)} \left( \begin{array}{c}
  0 \\
  0
\end{array} \right) = \left( \begin{array}{c}
  012 \\
  101
\end{array} \right)^0 = \left( \begin{array}{c}
  012 \\
  101
\end{array} \right) = \left( \begin{array}{c}
  0 \\
  1
\end{array} \right) \left( \begin{array}{c}
  2 \\
  1
\end{array} \right) \) and \( \text{val}_{(U,U)} \left( \begin{array}{c}
  102 \\
  001
\end{array} \right) = \left( \begin{array}{c}
  -2 \\
  4
\end{array} \right) \). Moreover, we have \( \text{val}_{(U,U)} \left( \begin{array}{c}
  101 \\
  011
\end{array} \right) = \left( \begin{array}{c}
  -2 \\
  1
\end{array} \right) \).
3 Multidimensional regular sequences

Definition 6. Let $\mathbf{U} = (U_1, \ldots, U_d)$ be a $d$-tuple of numeration systems. A sequence $f : \mathbb{N}^d \rightarrow \mathbb{K}$ is called $(\mathbf{U}, \mathbb{K})$-regular if the series $\sum_{n \in \mathbb{N}^d} f(n) \text{rep}_U(n)$ is $\mathbb{K}$-recognizable.

Example 7. Since the series defined by (1) is $\mathbb{N}$-recognizable, the sequence

$$f : \mathbb{N} \rightarrow \mathbb{N}, \ n \mapsto \begin{cases} 3|\text{rep}_U(n)|1 (1 + 3^{-g(\text{rep}_U(n))}) 2|\text{rep}_U(n)|2 & \text{if } g(\text{rep}_U(n)) \neq 0 \\ 3|\text{rep}_U(n)|1 \ 2|\text{rep}_U(n)|2 & \text{otherwise} \end{cases}$$

is $(U, \mathbb{N})$-regular.

The aim of this work is to prove the following result, which testifies to the robustness of the notion of $(\mathbf{U}, \mathbb{K})$-regular sequences.

Theorem 8. Suppose that $\mathbb{K}$ is commutative. For $f : \mathbb{N}^d \rightarrow \mathbb{K}$ and a $d$-tuple $\mathbf{U} = (U_1, \ldots, U_d)$ of Pisot numeration systems, the following assertions are equivalent.

1. The sequence $f$ is $(\mathbf{U}, \mathbb{K})$-regular.
2. For all finite alphabets $\mathbf{A} \subset \mathbb{Z}^d$, the series $\sum_{w \in \mathbf{A}^*} f(||\text{val}_U(w)||) w$ is $\mathbb{K}$-recognizable.
3. The series $\sum_{w \in \mathbf{A}^*_U} f(\text{val}_U(w)) w$ is $\mathbb{K}$-recognizable.
4. There exists a $\mathbb{K}$-recognizable series $S : \mathbf{A}^*_U \rightarrow \mathbb{K}$ such that for all $n \in \mathbb{N}^d$, $(S, \text{rep}_U(n)) = f(n)$.

4 Multidimensional normalization and delay

For a numeration system $U$, the normalization $\nu_U$ is the partial function from $\mathbb{Z}^*$ to $A^*_U$ which maps a word $w \in \mathbb{Z}^*$ with $\text{val}_U(w) \geq 0$ to the word $\text{rep}_U(\text{val}_U(w))$, see [13].

Theorem 9 ([13]). For any Pisot numeration system $U$ and any finite alphabet $\mathbf{A} \subset \mathbb{Z}$, there exists a computable DFA accepting the language $\{(v^U) \in (A \times A^U)^* : v \in 0^*\nu_U(u)\}$.

The previous result is usually referred to by saying that the normalization from the alphabet $A$ is computable by a finite automaton. A different construction of such a 2-tape automaton can be found in [6].

Example 10. The language $\{(v^U) \in \{0,1,2\}^2 : v \in 0^*\nu_U(u)\}$ is accepted by the DFA depicted in Figure 3.

Let us extend the definition of the normalization in order to obtain a total function. We keep the same vocabulary and notation. For a numeration system $U$, we let $\overline{A_U} = \{\overline{a} : a \in A_U\}$ and we define the (extended) normalization as the function

$$\nu_U : \mathbb{Z}^* \rightarrow (A_U \cup \overline{A_U})^*, \ w \mapsto \begin{cases} \text{rep}_U(\text{val}_U(w)) & \text{if } \text{val}_U(w) \geq 0 \\ \text{rep}_U(-\text{val}_U(w)) & \text{otherwise}. \end{cases}$$

Note that $A_U \cap \overline{A_U} = \{0\}$. Also note that $\nu_U(\mathbb{Z}^*) \subseteq A^*_U \cup (\overline{A_U})^*$.

Example 11. We have $\text{val}_U(\overline{22}) = -8$, so $\nu_U(\overline{22}) = \overline{\text{rep}_U(8)} = \overline{100}$. 

7
We generalize Theorem 9 to the normalization extended to negative values.

**Theorem 12.** For any Pisot numeration system $U$ and any finite alphabet $A \subset \mathbb{Z}$, there exists a computable DFA accepting the language $\{(u,v) \in (A \times (A_U^1 \cup A_U^-))^* : v \in 0^*\nu_U(u)\}$.

**Proof.** Let $A' \subset \mathbb{Z}$ be a symmetric alphabet of the form $\{-d, \ldots, d\}$ containing $A$. By Theorem 9, we can compute a DFA $N_1 = (Q, i, T, A', \delta_1)$ accepting the language $\{(u,v) \in (A' \times A_U^1)^* : v \in 0^*\nu_U(u)\}$. Define another DFA $N_2 = (Q, i, T, A' \times \overline{A_U^1}, \delta_2)$ where

$$\delta_2: Q \times (A' \times \overline{A_U^1}) \rightarrow Q, \ (q, (\frac{u}{v})) \mapsto \delta_1(q, (\frac{u}{v})).$$

By construction, $N_2$ accepts the language $\{(u,v) \in (A' \times \overline{A_U^1})^* : v \in 0^*\nu_U(u)\}$. Therefore, we can compute a third DFA $N'$ accepting the union of the previous two languages, which is equal to $\{(u,v) \in (A' \times (A_U^1 \cup \overline{A_U^1}))^* : v \in 0^*\nu_U(u)\}$. Finally, by deleting all transitions whose label has a first component not belonging to $A$, we obtain the desired DFA.  

The definition of the normalization can be generalized to the multidimensional setting by applying the normalization component-wise and padding with leading zeroes as defined in Section 2.3. For a $d$-tuple of numeration systems $U = (U_1, \ldots, U_d)$, we let $\overline{U} = \overline{A_{U_1}} \times \cdots \times \overline{A_{U_d}}$ and we define the (d-dimensional extended) normalization by

$$\nu_{U} : (\mathbb{Z}^d)^* \rightarrow (A_{U} \cup \overline{A_{U}})^*, \ \left(\begin{array}{c} w_1 \\ \vdots \\ w_d \end{array}\right) \mapsto \left(\begin{array}{c} \nu_{U_1}(w_1) \\ \vdots \\ \nu_{U_d}(w_d) \end{array}\right).$$

**Example 13.** We have $\nu_{(U,U)} \left(\frac{22}{10}\right) = \left(\frac{\nu_{U}(22)}{\nu_{U}(10)}\right)^0 = \left(\frac{100}{10}\right)^0 = \left(\frac{00}{01}\right)^0 = \left(\frac{1}{1}\right)(\frac{0}{0}) = \left(\frac{0}{0}\right)$. Note that $\nu_{(U,U)} \left(\frac{22}{10}\right) \notin A_{(U,U)}^* \cup (\overline{A_{U(U)}})^*$ since val$_{U}(22) < 0$ while val$_{U}(10) > 0$.

Let us define an operation on 2-tape automata in order to build multidimensional normalizers.
Definition 14. For each $j \in \{1, \ldots, d\}$, let $A_j = (Q_j, i_j, T_j, A_i \times B_i, \delta_j)$ be a 2-tape automaton. We define a $2d$-tape automaton $\bigotimes_{j=1}^d A_j = (Q, i, T, A \times B, \delta)$ by

- $Q = Q_1 \times \cdots \times Q_d$
- $i = (i_1, \ldots, i_d)$
- $T = T_1 \times \cdots \times T_d$
- $A = A_1 \times \cdots \times A_d$ and $B = B_1 \times \cdots \times B_d$
- $\delta: Q \times (A \times B) \to Q$, $((q_1, \ldots, q_d), (a, b)) \mapsto (\delta_1(q_1, (a_1^{b_1})), \ldots, \delta_d(q_d, (a_d^{b_d})))$ where $a = \left(\begin{array}{c} a_1 \\ \vdots \\ a_d \end{array}\right)$ and $b = \left(\begin{array}{c} b_1 \\ \vdots \\ b_d \end{array}\right)$.

The automaton $\bigotimes_{j=1}^d A_j$ can also be viewed as a 2-tape automaton where each component would be a $d$-tuple. We will make no distinction between these two equivalent points of view.

The following result generalizes Theorem 12 to the multidimensional setting.

Theorem 15. For any $d$-tuple of Pisot numeration systems $U = (U_1, \ldots, U_d)$ and any finite alphabet $A \subset \mathbb{Z}^d$, there exists a computable DFA accepting the language $\{(v) \in (A \times (A_U \cup \overline{A_U}))^*: v \in 0^*v_U(u)\}$.

Proof. Let $A_1, \ldots, A_d$ be finite alphabets of integers such that $A \subseteq A_1 \times \cdots \times A_d$. By Theorem 12 for each $j \in \{1, \ldots, d\}$, we can compute a DFA $N_j$ accepting the language $\{(v) \in (A_j \times (A_U \cup \overline{A_U}))^*: v \in 0^*v_U(u)\}$. Then the DFA $\bigotimes_{j=1}^d N_j$ accepts the language $\{(v) \in ((A_1 \times \cdots \times A_d) \times (A_U \cup \overline{A_U}))^*: v \in 0^*v_U(u)\}$. Indeed, let $(v) \in ((A_1 \times \cdots \times A_d) \times (A_U \cup \overline{A_U}))^*$ and write $u = \left(\begin{array}{c} u_1 \\ \vdots \\ u_d \end{array}\right)$ and $v = \left(\begin{array}{c} v_1 \\ \vdots \\ v_d \end{array}\right)$.

By definition, the 2$d$-tape automaton $\bigotimes_{j=1}^d N_j$ accepts $(v)$ if and only if for every $j \in \{1, \ldots, d\}$, $N_j$ accepts $(u)_j$. Finally, we obtain the desired DFA by removing from $\bigotimes_{j=1}^d N_j$ all transitions with a label $(a)_j$ whose first component $a$ does not belong to $A$. \qed

For a $d$-tuple of numeration systems $U = (U_1, \ldots, U_d)$ and a finite alphabet $A \subset \mathbb{Z}^d$ such that the language $\{(v) \in (A \times (A_U \cup \overline{A_U}))^*: v \in 0^*v_U(u)\}$ is regular, we define the (d-dimensional extended) normalizer $N_{U,A}$ to be the trim minimal automaton of this language. When the context is clear, we set $N_{U,A} = (Q_N, i_N, T_N, A \times (A_U \cup \overline{A_U}), \delta_N)$.

Moreover, when $A = A_U$ we simply write $N_U$ instead of $N_{U,A_U}$.

The normalizer $N_{U,A}$ is at the core of the reasoning leading to the proof of Theorem 8.

Let us make two useful observations. First, we note that the initial state is also final since the empty word is accepted by $N_{U,A}$. Second, the following lemma is a direct consequence of the unicity of normalized representations.
Lemma 16. For any two states \( q \) and \( q' \) of the normalizer \( \mathcal{N}_{U,A} \) and any letter \( a \in A \), there exists at most one letter \( b \in A_U \cup \bar{A}_U \) such that \( (q,(b^\ell),q') \) is a transition in \( \mathcal{N}_{U,A} \).

Example 21. Let \( \nu_U(w_1) = \nu_U(w_2) = \nu_U(w_3) = 100 \). Therefore, \( d_U(w_1) = 1, d_U(w_2) = 0 \) and \( d_U(w_3) = -1 \). For all \( \ell \in \mathbb{N} \), the normalizer \( \mathcal{N}_{U,A} \) accepts the words \( 0^{\ell+1}w_10^{\ell+1}0^{\ell+1}0^{\ell+1}z_1 \), \( 0^{\ell}w_20^{\ell}0^{\ell}0^{\ell}z_2 \) and \( 0^{\ell}w_30^{\ell}0^{\ell}0^{\ell}0^{\ell}z_3 \).

Now, let \( z_1 = 10, z_2 = 010 \) and \( z_3 = 0010 \). Then \( d_U(z_1) = 0, d_U(z_2) = -1 \) and \( d_U(z_3) = -2 \). In dimension 2, we have \( d_{U(U)}(w_1) = 1, d_{U(U)}(w_2) = 0, d_{U(U)}(w_3) = -1 \). For all \( \ell \in \mathbb{N} \), the normalizer \( \mathcal{N}_{U(U)} \) accepts the words \( 0^{\ell+1}w_10^{\ell+1}z_10^{\ell+1}0^{\ell+1}z_20^{\ell+1}0^{\ell+1}z_3 \).

5 Characterizations of \((U,K)\)-regular sequences

In this section, we consider a \( d \)-tuple of Pisot numeration systems \( U = (U_1, \ldots, U_d) \), a \( d \)-dimensional sequence \( f: \mathbb{N}^d \to K \) and a finite alphabet \( A \subset \mathbb{Z}^d \). In order to prove the implication \((1) \implies (2)\) of Theorem 6 some intermediate results and definitions are needed.

In what follows, we use the notation\(^1\)

\[
G_f = \sum_{n \in \mathbb{N}^d} f(n) \text{rep}_U(n) \quad \text{and} \quad V_{f,A} = \sum_{w \in A^*} f(||\text{val}_U(w)||)w.
\]

\(^1\)The letters \( G \) and \( V \) stand for "Greedy" and "Value" respectively.
If \( f \) is \((U, \mathcal{K})\)-regular, then by Proposition 2 there exists a \( \mathcal{K}\)-automaton recognizing the series \( G_f \). In this case, we let
\[
A_G = (Q_G, I_G, T_G, \mathbf{A}_U, E_G)
\]
be such a \( \mathcal{K}\)-automaton. Without loss of generality, by adding a new state if needed, we assume that \( A_G \) has a unique initial state \( i_G \), which has no incoming edge \([11, 17]\). By definition of \( G_f \), we have \((G_f, \mathbf{0}^k w) = 0\) for any \( k > 0 \) and any word \( w \in \mathbf{A}_U^* \). Therefore, we make the additional assumption that the unique initial state of \( A_G \) has no outgoing edge of label 0.

In order to construct a \( \mathcal{K}\)-automaton recognizing \( V_{f,A} \) starting from \( A_G \), we introduce an operation between a 2-tape DFA and a \( \mathcal{K}\)-automaton. For any two finite alphabets \( A, B \) and any coding \( \sigma: B^* \to B^* \), that is, a letter-to-letter morphism, this operation applied on a 2-tape automaton \( A \) over \( A \times B \) and a \( \mathcal{K}\)-automaton \( B \) over \( \sigma(B) \) produces a new \( \mathcal{K}\)-automaton \( A \odot B \) over \( A \). Roughly speaking, for all words \( u \in A^* \), if there exists a word \( v \in B^* \) such that \((u, v)\) is accepted by \( A \) then the new \( \mathcal{K}\)-automaton \( A \odot B \) mimics the behavior of the \( \mathcal{K}\)-automaton \( B \) on \( \sigma(v) \), that is, the weight of \( u \) in \( A \odot B \) is the weight of \( \sigma(v) \) in \( B \). Similar (but different) products of automata can be found in \([17]\).

**Definition 22.** Let \( A \) and \( B \) be two finite alphabets and let \( \sigma: B^* \to B^* \) be a coding. Let \( A = (Q_A, i_A, T_A, A \times B, \delta_A) \) be a DFA such that for any two states \( q, q' \in Q_A \) and any letter \( a \in A \), there exists at most one letter \( b \in B \) such that \( \delta_A(q, (b, a)) = q' \). Let \( B = (Q_B, I_B, T_B, \sigma(B), E_B) \) be a \( \mathcal{K}\)-automaton. With \( A \) and \( B \), we associate a new \( \mathcal{K}\)-automaton \( A \odot B = (Q, I, T, A, E) \) as follows.

- \( Q = Q_A \times Q_B \).
- \( I: Q \to \mathbb{K}, (q, q') \mapsto \begin{cases} I_B(q') & \text{if } q = i_A \\ 0 & \text{else.} \end{cases} \)
- \( T: Q \to \mathbb{K}, (q, q') \mapsto \begin{cases} T_B(q') & \text{if } q \in T_A \\ 0 & \text{else.} \end{cases} \)
- \( E: Q \times A \times Q \to \mathbb{K}, ((q_1, q_1'), a, (q_2, q_2')) \mapsto \begin{cases} E_B(q_1', \sigma(b), q_2') & \text{if } \exists b \in B, \delta_A(q_1, (b, a)) = q_2' \\ 0 & \text{else.} \end{cases} \)

Note that the extra assumption on \( A \) is required to ensure that \( E \) is a well-defined function. An illustration of this operation is given by the black part of Figure 5 in Example 25.

Consider the coding \( \sigma: (\mathbf{A}_U \cup \overline{\mathbf{A}}_U)^* \to \overline{\mathbf{A}}_U^*, w \mapsto ||w|| \). By Lemma 16, the \( \mathcal{K}\)-automaton \( N_{U,A} \odot A_G \) (with respect to the coding \( \sigma \)) is well defined. Let \( N_{U,A} \odot A_G = (Q, I, T, A, E) \). We now establish some properties of the latter \( \mathcal{K}\)-automaton. First, we note that \( N_{U,A} \odot A_G \) has a unique initial state, which has no incoming edge. Then we prove a technical lemma.

**Lemma 23.** For all \( q \in Q, a \in A \) and \( \ell_1, \ell_2 \in \mathbb{N} \), if there exist two paths labeled by \( \mathbf{0}^{\ell_1} a \) and \( \mathbf{0}^{\ell_2} a \) from \((i_N', i_G)\) to \( q \) in \( N_{U,A} \odot A_G \), then \( \ell_1 = \ell_2 \).

**Proof.** Let \((p, p')\) \( Q, a \in A \) and \( \ell_1, \ell_2 \in \mathbb{N} \) such that there exist two paths labeled by \( \mathbf{0}^{\ell_1} a \) and \( \mathbf{0}^{\ell_2} a \) from \((i_N', i_G)\) to \((p, p')\) in \( N_{U,A} \odot A_G \). Then there exist paths respectively
labeled by \((0^t a)^i\) and \((0^{t_2} a)^j\) for some \(v_1, v_2 \in (A_U \cup \overline{A_U})^*\) from \(i_N\) to \(p\) in \(N_{U,A}\). By hypothesis on \(A_C\), \(v_1\) and \(v_2\) cannot start with 0. Since \(N_{U,A}\) is co-accessible, there exists a path, say of label \((\delta^+)^i\), from \(p\) to a final state in \(N_{U,A}\). It follows that \(v_1 y = \nu_U(0^t ax) = \nu_U(0^{t_2} ax) = v_2 y\) which implies \(v_1 = v_2\), and in turn \(\ell_1 = \ell_2\) as desired. \(\square\)

Next, we study the series recognized by \(N_{U,A} \otimes A_C\).

**Lemma 24.** Let \(S\) be the series recognized by \(N_{U,A} \otimes A_C\). For all \(w \in A^*\), if \(d_U(w) \neq 0\) then the contribution of any path \(c \in C_{N_{U,A} \otimes A_C}(w)\) is zero, and

\[
(S, w) = \begin{cases} (G_f, ||\nu_U(w)||) & \text{if } d_U(w) = 0 \\ 0 & \text{otherwise.} \end{cases}
\]

**Proof.** Let \(w \in A^*\). If \(d_U(w) > 0\), then the contribution of any path in \(C_{N_{U,A} \otimes A_C}(w)\) is zero by Remark 20. Now, assume that \(d_U(w) < 0\). By Remark 20 the unique word \(v\) such that \(N_{U,A}\) accepts \((\delta^+)^i\) is \(v = 0^{-d_U(w)} \nu_U(w)\), which starts with 0. Since we assumed that there is no outgoing edge labeled by 0 from \(i_G\), there is no path labeled by \(w\) in \(N_{U,A} \otimes A_C\) starting from \((i_N, i_G)\). Therefore, if \(d_U(u) \neq 0\) then the contribution of any path in \(C_{N_{U,A} \otimes A_C}(w)\) is zero, and thus \((S, w) = 0\).

Suppose now that \(d_U(w) = 0\). Write \(w = a_1 \ldots a_t\) and \(\nu_U(w) = b_1 \ldots b_k\). There exists a unique sequence of states \(q_0, \ldots, q_{\ell}\) in \(Q_N\) such that \(q_0 = i_N, q_{\ell} \in T_N\) and for all \(i \in \{1, \ldots, \ell\}\), \(\delta_N((q_{i-1}, (a_i^+))) = q_i\). Let \(c = (c_0, a_1, q_1, q_1') \ldots (c_{\ell-1}, q_{\ell-1}, a_\ell, q_\ell)\) where \(c_0 = i_G\) and \(q_1', \ldots, q_\ell'\) are some states in \(Q_G\). Therefore,

\[
I(q_0, q_0') E(c) T(q_{\ell}, q_{\ell}') = I_G(i_G) \left( \prod_{i=1}^{\ell} E_G(q_{i-1}^{'}, ||b_i||, q_i') \right) T_G(q_{\ell}).
\]

We obtain

\[
(S, w) = \sum_{q_1', \ldots, q_{\ell}' \in Q_G} I_G(i_G) \left( \prod_{i=1}^{\ell} E_G(q_{i-1}^{'}, ||b_i||, q_i') \right) T_G(q_{\ell}) = (G_f, ||\nu_U(w)||).
\]

We now modify \(N_{U,A} \otimes A_C\) to create a new \(K\)-automaton \(A_{V,A}\), which will be proven to recognize the series \(V_{f,A}\). We define the \(K\)-automaton \(A_{V,A} = (Q_V, I_V, T_V, A, E_V)\) as follows.

- \(Q_V = Q \cup \{\alpha\}\).
- \(I_V: Q_V \to K\) is defined by \(I_V|_Q = I\) and \(I_V(\alpha) = I(i_N, i_G)\).
- \(T_V: Q_V \to K\) is defined by \(T_V|_Q = T\) and \(T_V(\alpha) = 0\).
- \(E_V: Q_V \times A \times Q_V \to K\) is defined as follows.
  
  \(E_V|_{(Q \times A \times Q)^3}(p, a, q) = \begin{cases} E(p, a, q) & \text{if } (p, a, q) \neq ((i_N, i_G), 0, (i_N, i_G)) \\ 1 & \text{else.} \end{cases}\)
2. For any \( q \in Q \) and \( a \in A \), \( E_V(\alpha, a, q) \) is equal to the sum of the weights \( E(c) \) of all paths \( c \) from \((i_N, i_G)\) to \( q \) in \( \mathcal{N}_{U,A} \circ \mathcal{A}_G \) labeled by \( 0^\ell a \) for some \( \ell \in \mathbb{N}_{\geq 1} \).

3. For all \( q \in Q_V \) and all \( a \in A \), \( E_V(q, a, \alpha) = 0 \).

Note that in Item 2, if such an integer \( \ell \) exists, then it is unique by Lemma 23, and moreover, it is bounded by the number of states in \( Q \).

The modification \( \mathcal{A}_{V,A} \) of \( \mathcal{N}_{U,A} \circ \mathcal{A}_G \) is symbolically depicted in Figure 4, where wavy arrows represent paths whereas straight arrows represent edges.

Let us emphasize that the new loop on the state \((i_N, i_G)\) was indeed not present in \( \mathcal{N}_{U,A} \circ \mathcal{A}_G \) since we took care to assume that the state \( i_G \) of the \( \mathbb{K} \)-automaton \( \mathcal{A}_G \) had no incoming edge. In view of Lemma 24, the \( \mathbb{K} \)-automaton \( \mathcal{N}_{U,A} \circ \mathcal{A}_G \) produces the correct coefficients of \( V_{f,A} \) only for words with zero delay. The role of the state \( \alpha \) is to deal with words in \( A^* \) with positive delays whereas the loop on the state \((i_N, i_G)\) is added to take into account words with negative delays. Note that \( \alpha \) might have outgoing edges of label \( 0 \). This happens whenever \( 0 \in A \) and there exist words \( w \in A^* \) having a delay greater than or equal to 2.

![Figure 4: The automaton \( \mathcal{A}_{V,A} \) is a modification of \( \mathcal{N}_{U,A} \circ \mathcal{A}_G \), which recognizes the series \( V_{f,A} \).](image.png)

We are now ready to prove our main result.

**Proof of Theorem 8.** The implications \( 2 \implies 3 \) and \( 3 \implies 4 \) are clear. For any \( \mathbb{K} \)-recognizable series \( S \) such that for all \( n \in \mathbb{N}^d \), \( (S, \text{rep}_{U}(n)) = f(n) \), we have \( G_f = S \circ \text{rep}_{U}(N^d) \). For every \( j \in \{1, \ldots, d\} \), \( \text{rep}_{U_j}(N^d) \) is a regular language since \( U_j \) is a Pisot numeration system. Therefore, the language \( \text{rep}_{U}(N^d) \) is regular as well. Then the characteristic series of \( \text{rep}_{U}(N^d) \) is \( \mathbb{K} \)-recognizable [14 Chapter 3, Proposition 2.1]. The implication \( 4 \implies 1 \) now follows from the fact that for a commutative semiring \( \mathbb{K} \), the Hadamard product of \( \mathbb{K} \)-recognizable series is \( \mathbb{K} \)-recognizable [18].
For all exactly two terms of the sum (3) are zero while the third one is equal to 1.

In order to obtain that $Cw$ we have to show that $C_{w} \in \mathbb{N}^{c} \subseteq A^{1}$. Let $w \in A^{*}$. Let $C_{1}$ be the set of those paths labeled by $w$ in $N_{w}^{c}$ going through the new loop on $(i_{N}, i_{G})$ at least once and let $C_{2}$ be the set of those paths labeled by $w$ in $N_{w}^{c}$ starting in the state $\alpha$. Note that each path in $C_{1}$ starts in the state $(i_{N}, i_{G})$ and that each path in $C_{A_{w}}(w) \setminus (C_{1} \cup C_{2})$ is a path in $N_{w}^{c} \otimes A_{G}$. Therefore, $(S', w)$ is equal to

$$
(S, w) + \sum_{c \in C_{1}} I_{V}(i_{N}, i_{G}) E_{V}(c) T_{V}(t_{c}) + \sum_{c \in C_{2}} I_{V}(\alpha) E_{V}(c) T_{V}(t_{c}).
$$

In order to obtain that $(S', w) = (V_{f, A}, w)$, we show that, depending on the delay $d_{U}(w)$, exactly two terms of the sum (3) are zero while the third one is equal to $(V_{f, A}, w)$.

Let us write $w = \left( \begin{array}{c} u_{1} \\ \vdots \\ u_{d} \end{array} \right)$.

For all $j \in \{1, \ldots, d\}$, since $\nu_{U}(w_{j}) \in A_{U}^{*} \cup (A_{U})^{*}$, we have that $\nu_{U}(\nu_{U}(w_{j})) = \nu_{U}(\nu_{U}(w_{j}))$. Then by definition of the series $G_{f}$ and $V_{f, A}$, we get that $(G_{f}, \nu_{U}(w))) = f(\nu_{U}(\nu_{U}(w))) = f(\nu_{U}(w))) = (V_{f, A}, w)$. It then follows from Lemma 24 that

$$
(S, w) = \begin{cases} 
(V_{f, A}, w) & \text{if } d_{U}(w) = 0 \\
0 & \text{else.}
\end{cases}
$$

Now, we consider the second term of (3). We show that

$$
\sum_{c \in C_{1}} I_{V}(i_{N}, i_{G}) E_{V}(c) T_{V}(t_{c}) = \begin{cases} 
(V_{f, A}, w) & \text{if } d_{U}(w) < 0 \\
0 & \text{else.}
\end{cases}
$$

For $c \in C_{1}$, let $\ell \in \mathbb{N}_{\geq 1}$ be the number of times $c$ goes through the loop on the initial state $(i_{N}, i_{G})$ and write $w = 0^l w'$. Then $d_{U}(w) = d_{U}(w') - \ell$. Let $c_{\ell}$ be the part of the path $c$ after reading $0^l$. Thus, $c_{\ell}$ is a path in $N_{w}^{c} \otimes A_{G}$ labeled by $w_{\ell}$ which starts in $(i_{N}, i_{G})$ and ends in $t_{c}$. If $d_{U}(w_{\ell}) \neq 0$, Lemma 24 implies that $I_{V}(i_{N}, i_{G}) E_{V}(c) T_{V}(t_{c}) = I(i_{N}, i_{G}) 1^{\ell} E(c_{\ell}) T(t_{c}) = 0$. Therefore, if $d_{U}(w) \geq 0$ then the contribution of any path $c \in C_{1}$ to the weight of $w$ is zero. Now suppose that $d_{U}(w) < 0$. Then only the paths $c \in C_{1}$ going exactly $-d_{U}(w)$ times through the loop on $(i_{N}, i_{G})$ may have a non-zero contribution. By setting $w' = w-d_{U}(w)$ and $c' = c-d_{U}(w)$, and by using (4), we get

$$
\sum_{c \in C_{1}} I_{V}(i_{N}, i_{G}) E_{V}(c) T_{V}(t_{c}) = \sum_{c' \in C_{N_{w}^{c} \otimes A_{G}}(w')} I(i_{N}, i_{G}) 1^{-d_{U}(w)} E(c') T(t_{c'}) = (S, w') = (V_{f, A}, w).
$$

We turn to the third term of (3). We show that

$$
\sum_{c \in C_{2}} I_{V}(\alpha) E_{V}(c) T_{V}(t_{c}) = \begin{cases} 
(V_{f, A}, w) & \text{if } d_{U}(w) > 0 \\
0 & \text{else.}
\end{cases}
$$
If \( w = \varepsilon \), the desired equality follows from the facts that \( d_U(\varepsilon) = 0 \) and \( T_V(\alpha) = 0 \). We now suppose that \( w \neq \varepsilon \) and write \( w = aw' \). Any path \( c \in C_2 \) can be decomposed as \( c = (\alpha, a, q)c' \) where \( q \in Q \) and \( c' \) is path in \( N_{U,A} \oplus A_G \) labeled by \( w' \). For each \( q \in Q \), we let \( P_q \) be the set of those paths in \( N_{U,A} \oplus A_G \) labeled by \( w' \) and starting in \( q \). By definition, the transition \((\alpha, a, q)\) has a non-zero weight in \( A_{V,A} \) if there exists a path in \( N_{U,A} \oplus A_G \) labeled by \( 0^a a \) for some \( \ell \in \mathbb{N} \). By Lemma 24 all such paths have the same label \( 0^a a \). Then

\[
\sum_{c \in C_2} I_V(\alpha)E_V(c)T_V(t_c) = \sum_{q \in Q} \sum_{c' \in P_q} I_V(\alpha)E_V(\alpha, a, q)E_V(c')T_V(t_{c'})
\]

\[
= \sum_{q \in Q} \sum_{c' \in P_q} I(i_N', i_G) \left( \sum_{p \in C_{U,A} \otimes A_G(0^a a)} E(p) \right) E(c')T(t_{c'})
\]

\[
= \sum_{q \in Q} \sum_{p \in C_{U,A} \otimes A_G(0^a a)} \sum_{c' \in P_q} I(i_N', i_G) E(p, c')T(t_{c'}). \tag{5}
\]

If \( \ell_{a,q} \neq d_U(\varepsilon) \), we have \( d_U(0^a a) = 0 \) and then it follows from Lemma 24 that each path \( p c' \) as in (5) has a zero contribution. In particular, if \( d_U(\varepsilon) \leq 0 \) then each term of the sum is equal to 0. Now, we suppose that \( d_U(\varepsilon) > 0 \). Then the sum (5) is equal to

\[
\sum_{c \in C_{U,A} \otimes A_G(0^a a)} I(i_N', i_G) E(c)T(t_c) = (S, 0^a a)_w.
\]

By (4) applied to \( 0^a a_w \), we have \( (S, 0^a a)_w = (V_{f,A}, 0^a a)_w = (V_{f,A}, w) \), hence the conclusion.

**Example 25.** With our notation, if \( f \) is the sequence (2), then the series \( G_f \) is equal to (1) and the series \( V_{f,\{0,1,2\}} \) is given by

\[
(V_{f,\{0,1,2\}}, w) = \begin{cases} 3^{|w_U(\varepsilon)|} 2^{|w_U(\varepsilon)|} + 3^{|w_U(\varepsilon)|} g(w_U(\varepsilon)) 2^{l_U(\varepsilon)} & \text{if } g(w_U(\varepsilon)) \neq 0 \\ 3^{l_U(\varepsilon)} 2^{l_U(\varepsilon)} & \text{if } g(w_U(\varepsilon)) = 0. \end{cases}
\]

for all \( w \in \{0,1,2\}^* \). Let us call \( A_G \) and \( A_{V,\{0,1,2\}} \) the \( K \)-automata of Figures 1 and 5 respectively. Note that \( A_G \) satisfies our assumptions: it recognizes \( G_f \) and has a unique initial state, which has no incoming edge and no outgoing edge of label 0. By looking at the normalizer \( N_U \) in Figure 3 we see that for all \( w \in \{0,1,2\}^* \), the delay \( d_U(w) \) is at most 1. This implies that in \( A_{V,\{0,1,2\}} \), the state \( a \) has no outgoing edge labeled by 0. By our construction, the series \( V_{f,\{0,1,2\}} \) is recognized by \( A_{V,\{0,1,2\}} \), which is precisely the modified version of \( N_U \oplus A_G \).

We end this section by making two points.

**Remark 26.** The commutativity of the semiring \( K \) is only needed to prove the implication (4) \( \implies \) (1). Indeed, this hypothesis is required to ensure that the family of \( K \)-recognizable series is closed under the Hadamard product.
Remark 27. In the proof of Theorem 8, the Pisot hypothesis might be too strong. For proving the implication (1) \(\Rightarrow\) (2), it is merely used in order to get that the normalization is computable by a finite automaton. For the implication (4) \(\Rightarrow\) (1), we only need the assumption that the numeration systems \(U_j\) all have regular numeration languages \(\text{rep}_{U_j}(\mathbb{N})\). The other implications do not need any assumption on the numeration systems.

6 From a greedy-based linear representation to a value-based one

The proof of Theorem 8 is constructive. In particular, by using the correspondence between \(K\)-automata and linear representations, we provide an algorithm that, starting with a linear representation of the series \(G_f = \sum_{n \in \mathbb{N}^d} f(n) \text{rep}_U(n)\) and a finite alphabet \(A \subset \mathbb{Z}^d\), computes a linear representation of the series \(V_{f,A} = \sum_{w \in A^*} f(||\text{val}_U(w)||) w\).

Let \((\lambda, \mu, \gamma)\) be a linear representation of \(G_f\) with \(\lambda \in K^{1 \times r}\), \(\mu: A_U \rightarrow K^{r \times r}\) and \(\gamma \in K^{r \times 1}\) corresponding to the \(K\)-automaton \(A_G\). In particular, \(r\) is the number of states of \(A_G\) and for all \(i, j \in \{1, \ldots, r\}\) and all \(a \in A_U\), \(\mu(a)_{i1} = 0, \mu(0)_{ij} = 0\) and \(\lambda_i = 0\) if \(i \neq 1\).

Using Definition 22, we first build a linear representation \((L, M, R)\) of the series recognized by \(N_{U,A} \otimes A_G\) with \(L \in K^{1 \times rs}\), \(M: A^* \rightarrow K^{rs \times rs}\) and \(R \in K^{rs \times 1}\). Let \(Q_N = \{q_1, \ldots, q_s\}\)
where \( q_1 \) is the unique initial state. We set
\[
L = (\lambda_1 \ 0 \ \cdots \ 0) \quad \text{and} \quad R = \left(\begin{array}{c} R_1 \\ \vdots \\ R_s \end{array}\right)
\]
where \( R_1, \ldots, R_s \in \mathbb{K}^{r \times 1} \) are given by
\[
R_i = \begin{cases} 
\gamma & \text{if } q_i \in T_N \\
0_{r \times 1} & \text{otherwise}
\end{cases}
\]
for all \( i \in \{1, \ldots, s\} \). For all \( a \in A \) and all \( i, j \in \{1, \ldots, s\} \), we define matrices \( M^{(a)}_{i,j} \in \mathbb{K}^{r \times r} \) by
\[
M^{(a)}_{i,j} = \begin{cases} 
\mu(||b||) & \text{if } \exists b \in A_U \cup \overline{A_U}, \delta_N(q_i, (b^a)) = q_j \\
0_{r \times r} & \text{otherwise}.
\end{cases}
\]
Finally, for all \( a \in A \), we set
\[
M(a) = \left(\begin{array}{c} M^{(a)}_{1,1} \\ \vdots \\ M^{(a)}_{s,1} \end{array} \cdots \begin{array}{c} M^{(a)}_{s,1} \\ \vdots \\ M^{(a)}_{s,s} \end{array}\right).
\]

The last step is to encode the modification \( A_{V,A} \) of \( \mathcal{N}_{U,A} \otimes \mathcal{A}_G \) that consists in adding the state \( a \) together with its transitions and the loop of label 0 on the initial state. A linear representation of the series \( V_{f,A} \) is given by \( (\lambda', \mu', \gamma') \) with \( \lambda' \in \mathbb{K}^{1 \times (rs+1)}, \mu' : A^* \to \mathbb{K}^{(rs+1) \times (rs+1)} \) and \( \gamma' \in \mathbb{K}^{(rs+1) \times 1} \) defined as follows. We set \( \lambda' = (\lambda_1 \ \ell) \) and \( \gamma' = \left(\begin{array}{c} 0_r \\ \ell \end{array}\right) \).

For each \( a \in A \) and each \( j \in \{1, \ldots, rs\} \), we let \( \ell_{a,j} \) be the unique positive \( \ell \) such that there exists a path labeled by \( 0^\ell a \) from the initial state to the \( j \)-th state of \( \mathcal{N}_{U,A} \otimes \mathcal{A}_G \) if it exists (see Lemma 23 and Figure 4). Here, we consider the order on the states of \( \mathcal{N}_{U,A} \otimes \mathcal{A}_G \) induced by the matrix indexing. For \( a \in A \), we define
\[
\mu'(a) = \left(\begin{array}{c} 0 \\ m(a) \end{array}\right)
\]
where \( M'(a) = M(a) \) if \( a \neq 0 \) and for all \( i, j \in \{1, \ldots, rs\} \),
\[
(M'(0))_{i,j} = \begin{cases} 1 & \text{if } i = j = 1 \\
(M(0))_{i,j} & \text{otherwise}
\end{cases}
\]
and
\[
m(a)_j = \begin{cases} M(0^\ell a)_{1,j} & \text{if } \ell_{a,j} \text{ exists} \\
0 & \text{otherwise}.
\end{cases}
\]

To compute \( m(a) \) for each \( a \in A \), we can either look into the \( \mathbb{K} \)-automaton \( \mathcal{N}_{U,A} \otimes \mathcal{A}_G \) in order to determine the exponents \( \ell_{a,j} \) for \( j \in \{1, \ldots, rs\} \), or we can use the observation that
\[
m(a) = \left(\begin{array}{c} \sum_{\ell=1}^{rs} M(0^\ell a)_{1,1} \\ \vdots \\ \sum_{\ell=1}^{rs} M(0^\ell a)_{1,rs} \end{array}\right)
\]
since for all \( j \in \{1, \ldots, rs\} \), \( \ell_{a,j} \leq rs \) and \( M(0^\ell a)_{1,j} = 0 \) if \( \ell \neq \ell_{a,j} \).

In practice, removing non-accessible and non-coaccessible states often drastically reduce the size of the linear representation of \( V_{f,A} \).
Example 28. The linear representation $(\lambda', \mu', \gamma')$ corresponding to the automaton $A_{V,\{0,1,2\}}$ of Figure [2] with respect to the order of the states

$\alpha, (1, W), (1, X), (1, Z), (2, Y), (3, X), (3, Y), (3, Z), (4, X), (5, X)$

is given by $\lambda' = (1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0)$, $\gamma' = (0 \ 1 \ 1 \ 1 \ 1 \ 0 \ 0 \ 0 \ 0 \ 1)^T$ and

$\mu'(0) = \begin{pmatrix}
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 3 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
\end{pmatrix}$, $\mu'(1) = \begin{pmatrix}
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 3 & 1 & 0 & 2 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
\end{pmatrix}$, $\mu'(2) = \begin{pmatrix}
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
\end{pmatrix}$.

7 Acknowledgment

Célia Cisternino is supported by the FNRS Research Fellow grant 1.A.564.19F. Manon Stipulanti is supported by a Francqui Foundation Fellowship of the Belgian American Educational Foundation.

References

[1] Jean-Paul Allouche, Klaus Scheicher, and Robert F. Tichy. Regular maps in generalized number systems. Math. Slovaca, 50(1):41–58, 2000.

[2] Jean-Paul Allouche and Jeffrey Shallit. The ring of $k$-regular sequences. Theoret. Comput. Sci., 98(2):163–197, 1992.

[3] Jean-Paul Allouche and Jeffrey Shallit. Automatic sequences. Cambridge University Press, Cambridge, 2003. Theory, applications, generalizations.

[4] Jean Berstel and Christophe Reutenauer. Noncommutative rational series with applications, volume 137 of Encyclopedia of Mathematics and its Applications. Cambridge University Press, Cambridge, 2011.

[5] Anne Bertrand-Mathis. Comment écrire les nombres entiers dans une base qui n’est pas entière. Acta Math. Hungar., 54(3-4):237–241, 1989.

[6] Véronique Bruyère and Georges Hansel. Bertrand numeration systems and recognizability. volume 181, pages 17–43. 1997. Latin American Theoretical INformatics (Valparaíso, 1995).

[7] J. Richard Büchi. Weak second-order arithmetic and finite automata. Z. Math. Logik Grundlagen Math., 6:66–92, 1960.

[8] Émilie Charlier. First-order logic and numeration systems. In Sequences, groups, and number theory, Trends Math., pages 89–141. Birkhäuser/Springer, Cham, 2018.

[9] Émilie Charlier, Narad Rampersad, and Jeffrey Shallit. Enumeration and decidable properties of automatic sequences. Internat. J. Found. Comput. Sci., 23(5):1035–1066, 2012.
[10] Alan Cobham. On the base-dependence of sets of numbers recognizable by finite automata. *Math. Systems Theory*, 3:186–192, 1969.

[11] Samuel Eilenberg. *Automata, languages, and machines. Vol. A*. Academic Press [A subsidiary of Harcourt Brace Jovanovich, Publishers], New York, 1974. Pure and Applied Mathematics, Vol. 58.

[12] Christiane Frougny and Jacques Sakarovitch. Number representation and finite automata. In *Combinatorics, automata and number theory*, volume 135 of *Encyclopedia Math. Appl.*, pages 34–107. Cambridge Univ. Press, Cambridge, 2010.

[13] Christiane Frougny and Boris Solomyak. On representation of integers in linear numeration systems. In *Ergodic theory of $\mathbb{Z}^d$ actions (Warwick, 1993–1994)*, volume 228 of *London Math. Soc. Lecture Note Ser.*, pages 345–368. Cambridge Univ. Press, Cambridge, 1996.

[14] Julien Leroy, Michel Rigo, and Manon Stipulanti. Counting the number of non-zero coefficients in rows of generalized Pascal triangles. *Discrete Math.*, 340(5):862–881, 2017.

[15] M. Lothaire. *Algebraic combinatorics on words*, volume 90 of *Encyclopedia of Mathematics and its Applications*. Cambridge University Press, Cambridge, 2002.

[16] Michel Rigo and Arnaud Maes. More on generalized automatic sequences. *J. Autom. Lang. Comb.*, 7(3):351–376, 2002.

[17] Jacques Sakarovitch. *Elements of Automata Theory*. Cambridge University Press, Cambridge, 2009. Translated from the 2003 French original by Reuben Thomas.

[18] Marcel Paul Schützenberger. On a theorem of R. Jungen. *Proc. Amer. Math. Soc.*, 13:885–890, 1962.

[19] Édouard Zeckendorf. Représentation des nombres naturels par une somme des nombres de Fibonacci ou de nombres de Lucas. *Bull. Soc. Roy. Sci. Liège*, 41:179–182, 1972.