A Secure Dynamic Job Scheduling on Smart Grid using RSA Algorithm

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Abstract — Grid computing is a computation methodology using group of clusters connected over high-speed networks that involves coordinating and sharing computational power, data storage and network resources. Integrating a set of clusters of workstations into one large computing environment can improve the availability of computing power. The goal of scheduling is to achieve highest possible system throughput and to match the application need with the available computing resources. A secure scheduling model is presented, that performs job grouping activity at runtime. In a Grid environment, security is necessary because grid is a dynamic environment and participates are independent bodies with different policies, objectives and requirements. Authentication should be verified for Grid resource owners as well as resource requesters before they are allowed to join in scheduling activities. In order to achieve secure resource and job scheduling including minimum processing time and maximum resource utilization, A Secure Resource by using RSA algorithm on Networking and Job Scheduling model with Job Grouping strategy (JGS) in Grid Computing has been proposed. The result shows significant improvement in the processing time of jobs and resource utilization as compared to dynamic job grouping (DJG) based scheduling on smart grids (SG).

Keywords — RSA, SMART GRID, DYNAMIC JOB GROUPING, JOB GROUPING STRATEGY.

1. INTRODUCTION:

Wellner defines grid technology as “the technology that enables resource virtualization, on-demand provisioning, and service (resource) sharing between organizations.” IBM defines grid computing as “the ability, using a set of open standards and protocols, to gain access to applications and data, processing power, storage capacity and a vast array of other computing resources over the Internet. A grid is a type of parallel and distributed system that enables the sharing, selection, and aggregation of resources distributed across ‘multiple’ administrative domains based on their (resources) availability, capacity, performance, cost and users' quality-of-service requirements.” In 1965 Fernando Corbató and the other designers of the Multics operating system envisioned a computer facility operating “like a power company or water company”. Buyya/Venugopal define grid as “a type of parallel and distributed system that enables the sharing, selection, and aggregation of geographically distributed autonomous resources dynamically at runtime depending on their availability, capability, performance, cost, and users' quality-of-service requirements”. CERN, one of the largest users of grid technology, “a service for sharing computer power and data storage capacity over the Internet.” Grid computing is designed to work independent problems in parallel, thereby leveraging the computer processing power of a distributed model. Prior to grid computing, any advanced algorithmic process was only available with super computers. Grid computing is a term referring to the combination of computer resources from multiple administrative domains to reach a common goal.

Ian Foster lists these primary attributes:
(i) Computing resources are not administered centrally.
(ii) Open standards are used.

GRID ARCHITECTURE:

Figure 1.1 Layered Grid Architecture

In figure 1.1 Grid architecture defines a layered collection of protocols.
1. Fabric layer
2. Connectivity layer
3. Resource layer
4. Collective layer

(1) Fabric layer includes the protocols and interfaces that provide access to the resources that are being shared, including computers, storage systems, datasets, programs, and networks. This layer is a logical view rather than a physical view. For example, the view of a cluster with a local resource manager is defined by the local resource manager, and not the cluster hardware. Likewise, the fabric provided by a storage
system is defined by the file system that is available on that system, not the raw disk or tapes.

(2) The connectivity layer defines core protocols required for Grid-specific network transactions. This layer includes the IP protocol stack (system level application protocols transport and internet layers), as well as core Grid security protocols for authentication and authorization.

(3) Resource layer defines protocols to initiate and control sharing of (local) resources. Services defined at this level are gatekeeper, GRIS, along with some user oriented application protocols from the Internet protocol suite, such as file-transfer.

(4) Collective layer defines protocols that provide system oriented capabilities that are expected to be wide scale in deployment and generic in function. This includes bandwidth brokers, resource brokers.

**Resource management in grid computing** For running applications, resource management and job scheduling are the most crucial problems in grid computing systems. In future the development of grid technology, it is very likely that corporations, universities and public institutions will exploit grids to enhance their computing infrastructure. The goal of scheduling is to achieve highest possible system throughput and to match the application need with the available computing resources.

In figure 1.2 basic grid model generally composed of a number of hosts, each composed of several computational resources, which may be homogeneous or heterogeneous.

The four basic building blocks of grid model are user, resource broker, grid information service (GIS) and Lastly resources. When user requires high speed execution, the job is submitted to the broker in grid. Broker splits the job into various tasks and distributes to several resources according to user’s requirements and availability of resources. GIS keeps the status information of all resources which helps the broker for scheduling.

**GridSim:**

GridSim is a Java-based toolkit for modeling, and simulation of distributed resource management and scheduling for conventional Grid environment. GridSim is based on SimJava, a general purpose discrete-event simulation package implemented in Java. All components in GridSim communicate with each other through message passing operations defined by SimJava.

**Features of GridSim:**

1. It allows modeling of heterogeneous types of resources
2. Resources can be modeled operating under space or time shared mode
3. Resource capability can be defined in the form of MIPS (Million Instructions Per Second).
4. Resources can be located in any time zone.
5. Resources can be booked for advance reservation.
6. Applications with different parallel application models can be simulated.

GridSim toolkit is suitable for application scheduling simulations in Grid Computing environment. GridSim is of great value to both students and experienced researchers who want to study Grids, or test new algorithms and strategies in a controlled environment. By using GridSim, they are able to perform repeatable experiments and studies that are not possible in a real dynamic Grid environment.

**Advantages of using GridSim are:**

1. Various allocation or scheduling policies can be made and integrated into GridSim easily, by extending them from one of the classes.
2. Has the infrastructure or framework to support advance reservation, auction and Grid functionalities of a Grid system.
3. Has the ability that reads workload traces taken from supercomputers for simulating a realistic Grid environment? This functionality is useful for testing a resource scheduling problem. This is useful for simulating over a public network where the network is congested.

Research students in the GRIDS Laboratory are themselves heavy users of GridSim and extend it whenever
necessary for their own research needs. In the last 5 years, GridSim has been continuously extended in this manner to include many new capabilities and has also received contributions from external collaborators.

The Job scheduler performs the job grouping and resource selection process. Once job groups are created, they are immediately sent for computation to the selected clusters accordingly and lastly, results are sent back to the respective Grid users.

**ISSUES:**

The issue of the secure resource and job scheduling model with job grouping strategy is to overcome the deficiencies of dynamic job grouping strategy. The deficiencies are:

(a) There is no security in dynamic job grouping based model

(b) The processing time to compute jobs is high in dynamic grouping based model.

In Secure resource and job scheduling model with job grouping strategy

(a) The registration process is proposed and authentication is verified for the resource requesters to join in scheduling activity

(b) Secure resource and job scheduling model with job grouping strategy contains all the Information about the resources of clusters such as MIPS (Million instructions per Second), bandwidth (in MBPS) and file size (Mb) in the global scheduler and assigns the resources to the jobs by considering all the constraints of resources. Hence the processing time of jobs will be decreased.

2. **JOB SCHEDULING STRATEGIES**

The Chapter Job scheduling strategies deals with the Researchers view of grid computing, scheduling objectives and hierarchal Job Scheduling for Cluster of Workstations

**2.1 RESEARCHERS VIEW OF GRID COMPUTING**

Grid computing is a term referring to the combination of computer resources from multiple administrative domains to reach a common goal. Foster, 2003 stated that the concept of forming pool of resources has led the way into creating large computing facilities which are currently popular known as Grid computing.

As soon as computers are interconnected and communicating, due to the faster and more capable hardware and network components and increasingly sophisticated software led to effective and efficient utilization of widely distributed resources to fulfill a range of application needs.

Many researchers are working in the field of wide-area distributed computing. These research groups have implemented middleware; libraries and tools that allow the cooperative use of geographically distributed resources to act as a single powerful platform for the execution of a range of parallel and distributed applications.
This approach to computing has been known by several names, such as metacomputing, scalable computing, global computing, Internet computing and lately known as grid computing.

Grids are often constructed with the aid of general-purpose grid software libraries known as middleware. One of the main strategies of grid computing is to use middleware to divide and apportion pieces of a program among several computers, sometimes up to many thousands. Middleware is generally considered to be the layer of software sandwiched between the operating system and applications, providing a variety of services required by an application to function correctly.

According to F.Berman middleware has re-emerged as a means of integrating software applications running in distributed heterogeneous environments. In a Grid, the middleware is used to hide the heterogeneous nature and provide users and applications with a homogeneous and seamless environment by providing a set of standardized interfaces.

2.2 Scheduling Objectives

Job scheduling means determining when and where the jobs are executed and how many resources are allocated. W.E.Welli stated that the basic objectives for job schedulers are to minimize response time and maximize utilization although typically there is some trade-off between both.

The goal of scheduling is to achieve the highest possible system throughput and to match the application need with the available computing resources. Scheduling jobs in a Grid Computing environment is a critical problem, as it involves the usage of heterogeneous resources that are geographically distributed. Job scheduling is the process of mapping jobs to the specific available physical resources and tries to minimize the processing time of the jobs. Hence there is a need for scheduling algorithms for efficient scheduling.

According to 2010 2nd International Conference on Signal Processing Systems the scheduling algorithms are defined as:

A. Efficient Utilization of Computing Resources Using Highest Response Next Scheduling in Grid (HRN)
B. Node Allocation in Grid Computing Using Optimal Resource Constraint (ORC) Scheduling
C. Hierarchical Job Scheduling for Clusters of Workstations (HJS)
D. Resource Co Allocation for Scheduling Tasks with Dependencies in grid (RCSTD)
E. Job Schedule Model Based on grid environment (JSMB)
F. Dynamic Job Grouping-Based Scheduling for Deploying Applications with Fine Grained Tasks on Global Grids (DJGBS)
G. Scheduling Framework for Bandwidth-Aware Job Grouping-Based scheduling in grid computing (SFBAJG)
H. A Bandwidth-Aware Job Grouping-Based Scheduling on Grid Environment (BAJGBS)

Job Grouping Strategy

The World Academy of Science, Engineering and Technology 64 2010 describes this grouping strategy maximizes the utilization of Grid resources, reduces processing time of jobs and network delay to schedule and execute jobs on the Grid. Grouping strategy is based on processing capability (in MIPS), bandwidth (in Mb/s), and memory-size (in Mb) of the available resources.

Jobs are grouped according to the capability of the selected resource. Therefore, the following conditions must be satisfied:

(i) GroupedJob_MI ≤ Resource_MIPS * Granularity size
(ii) GroupedJob_MS ≤ Resource_MS
(iii) GroupedJob_MS ≤ Resource_baud_rate * Tcomm

Where, MI (Million Instruction) is job’s required computational power, MIPS (Million Instruction Per Second) is processing capability of the resource. Granularity size is user defined time which is used to measure total number of jobs that can be completed within a specified time, Grouped job_MS is required Memory Size of group job, Resource_MS is the amount of Memory available at resource Baud rate is the bandwidth capacity of resource, and Tcomm is the job’s communication time Eq (i) required computational power of grouped job shouldn’t exceed to the resource’s processing capability. Eq (ii) Memory-size requirement of grouped job shouldn’t exceed to the resource’s memory-size capability. Eq (iii) Memory-size of the grouped job shouldn’t exceed to resource’s transfer capability within a specific time period.

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information table and sends the acquired information to the grouping and selection services.

After getting information the grouping and selecting services sends the grouped jobs to the dispatcher. Dispatcher dispatches the grouped jobs and finds the resources and the jobs are computed.

Hierarchical Job Scheduling for Clusters of Workstations

According to the Proceedings of the sixth annual conference of the Advanced School for Computing and Imaging, stated that advent of large computing power in workstations and high speed networks, the high performance computing community is moving from the use of massively parallel processors (MPPs) to cost effective clusters of workstations. Integrating a set of clusters of workstations into one large computing environment can improve the availability of computing power, e.g. Globus. The scheduling system for large environments can be divided into two levels scheduling across clusters (wide area scheduling) and scheduling within a cluster. Zhou et al introduced two level time-sharing scheduling for parallel and sequential jobs to attain a good performance for parallel jobs and a short turn-around time for sequential jobs.

At the upper level, time slots are used. Each time slot is divided into 2 slots (one for sequential jobs and a second for tasks of parallel jobs). However at the local level a task of a parallel job may share its time slot with sequential jobs. K.Y Wang uses hierarchical decision scheduling that is implemented using global and local schedulers. The global scheduler is responsible for long term allocation of system resources, and the local scheduler is responsible for short term decisions concerning processor allocation.

Implementation of Secure Resource RSA Algorithm and Job Scheduling Algorithm on Smart Grid

Implementation deals with Secure Resource and Job scheduling Architecture, modules and System design, User and Resource registration Algorithm, Job Grouping Algorithm, Ellipse and SQLserver2008 installations. The Secure Resource and Job scheduling model is divided into three levels

(1) user level
(2) global level and
(3) cluster level.
from the Grid information system (GIS) and sends it back to the user.

**Sequence diagram for secure resource and job scheduling model with job grouping**

Sequence diagrams are used to represent or model the flow of messages, events and actions between the objects or components of a system. Sequence Diagrams are used primarily to design, document and validate the architecture, interfaces and logic of the system by describing the sequence of actions that need to be performed to complete a task. In the figure 3.3 the resources need to be registered to Grid information service. Grid user submits jobs to Admin. Admin checks for the availability of resources and requests the available resource information from the Grid information system. The resource list is provided to Admin through Grid information service. The Admin request the resource characteristics from Grid information system and gets the resource characteristics list. Admin submits the grouped jobs and result is send to the user.

**User resource registration algorithm**

**Input:** Username and register resource

**Output:** authenticates the user and checks the availability of the resources

**Step1**: User name and user id is authenticated

var : user, resource
for (i = 1; i<Nu; i++)
{Nu = Number of users.}//
getrequest(specification);

**Sequence diagram for dynamic job grouping based scheduling**

In the figure 3.4 the resources need to be registered to Grid information service. Grid user submits jobs to Admin. Admin checks for the availability of resources and requests the available resource information from the Grid information system.

The resource list is provided to Admin through Grid information service. The Admin request the resource characteristics from Grid information system and gets the resource characteristics list.

Admin submits the grouped jobs, if grouped job (MI) is greater than resource (MI) send that job to joblist1. The job in joblist1 should be grouped and the Joblist1 needs to wait until the jobs in the joblist assigned with resources.

**User resource registration algorithm**

**Input:** Username and register resource

**Output:** authenticates the user and checks the availability of the resources

**Step1**: User name and user id is authenticated

var : user, resource
for (i = 1; i<Nu; i++)
{Nu = Number of users.}//
getrequest(specification);
Step2: Resource name, Resource enter time, Resource MIPS, Resource bandwidth, Resource memory size are registered to database for (i = 1; i<Nr; i++)

Step3: Scheduler checks the availability of resources Submit job to scheduler ()

Step4: If resources are available executes the job grouping algorithm
call job grouping algo();
scheduler submits job to resource;
resource computed jobs;
end back result to users;

Job grouping algorithm
Input : JobLength
Output : Resources and clusters
Step1 : Submits the user jobs to the grid
Step2 : Receive job list, JList
Step3 : Receive resource list, Rlist
for i:= 0 to JobList Size-1 do
for j:=O to Rlist_Size-1 do

Step4 : Grouped Job(million instructions) should not be equal to zero
GroupedjobMI:= 0
Step5 : Resource Million instructions can be obtained by multiplying resource MIPS
(processing capability) with granularity size which is taken as 3 in the project
R(mi) := Rlist_MIPS* Granularity Size
Step6: If Grouped job million instructions is less to Resource million instructions
Add another job
while GroupedjobMI < R(mi)J and i < JobList_Size-1 do

Groupedjob(MI):=Groupedjob(MI)+joblisti_(MI)
i++
Step7: If Grouped job million instructions is greater than Resource million instructions
if GroupedjobMI > R(mi)J then

Step8: The loop is repeated between the available resources and should assign the resource.
Step9: Place the Grouped job(j) to Target Resource List(j) for computation
Step10: Receive computed Grouped job from Resource List j.

RSA Algorithm for Smart Grids Security:
RSAParameters m_private;
RSAParameters m_public;
public byte[] GetHash(byte[] message)
{
   MD5CryptoServiceProvider md5 = new MD5CryptoServiceProvider();
   return md5.ComputeHash(message, 0, message.Length);
}

public byte[] CreateSignature(byte[] hash)
{
   RSACryptoServiceProvider RSA = new RSACryptoServiceProvider();
   RSAPKCS1SignatureFormatterRSAFormatter = new RSAPKCS1SignatureFormatter(RSA);
   RSAFormatter.SetHashAlgorithm("MD5");
   m_public = RSA.ExportParameters(false);
   m_private = RSA.ExportParameters(true);
   return RSAFormatter.CreateSignature(hash);
}

public bool VerifySignature(byte[] hash, byte[] signedhash)
{
   RSACryptoServiceProvider RSA = new RSACryptoServiceProvider();
   RSAParameters RSAKeyInfo = new RSAParameters();
   RSAKeyInfo.Modulus = m_public.Modulus;
   RSAKeyInfo.Exponent = m_public.Exponent;
   RSA.ImportParameters(RSAKeyInfo);
   RSAPKCS1SignatureDeformatter RSADeformatter = new RSAPKCS1SignatureDeformatter(RSA);
   RSADeformatter.SetHashAlgorithm("MD5");
   return RSADeformatter.VerifySignature(hash, signedhash);
}

Experimental results
In the Chapter "Secure resource and a job scheduling model with a job grouping strategy has to provide security for the grid information system."
Admin or the server has to validate the user id and password of the user.

A new user must register before it should enter into scheduling process.

After the registration process the user can submit jobs, view profile and view assigned resources.

The admin can view profiles, view user jobs, View resources, Assign resources, view assigned resources, assign cluster. View cluster and finds the processing time of secure resource and job scheduling algorithm and dynamic job grouping algorithm.

The profile that is entered by a new user will be updated in the Admin module.

Resources MIPS, resource bandwidth, name of the resource and resource enter time and resource memory size are registered to database.

The admin will also provide the cluster to process the group of jobs.

To process these jobs the Secure resource and a job scheduling model (SRJM) takes less time comparing to dynamic job grouping based algorithm.(DJGB).

In the module of SRJM, contains the block of SRJM and DJGB.

SRJM is secured and it can be accessed by the authorized user.

The processing time of jobs in SRJM is less compared to DJGB.

The SRJM algorithm is written in java in the project.

A simulation is conducted in heterogeneous environment where each cluster has machines with different characteristics and MIPS, to verify the improvement of proposed model over other dynamic job grouping based scheduling algorithm.

In this simulation, size of the granularity is taken as 3 seconds for both scheduling algorithm.

| Resource Name | Resource MIPS | Resource Bandwidth | Resource Memory Size | Granularity Size |
|---------------|---------------|---------------------|----------------------|------------------|
| R1            | 10            | 100                 | 100                  | 3                |
| R2            | 20            | 150                 | 200                  | 3                |
| R3            | 30            | 200                 | 300                  | 3                |
| R4            | 40            | 250                 | 400                  | 3                |

| Table registered resources (R1 to R4) |

| Resource Name | Resource MIPS | Resource Bandwidth | Resource Memory Size | Granularity Size |
|---------------|---------------|---------------------|----------------------|------------------|
| R5            | 50            | 300                 | 500                  | 3                |
| R6            | 60            | 350                 | 600                  | 3                |
| R7            | 70            | 400                 | 700                  | 3                |
| R8            | 80            | 450                 | 800                  | 3                |

| Table registered resources (R4 to R8) |

| Resource Name | Resource MIPS | Resource Bandwidth | Resource Memory Size | Granularity Size |
|---------------|---------------|---------------------|----------------------|------------------|
| R9            | 90            | 500                 | 900                  | 3                |
| R10           | 100           | 550                 | 1000                 | 3                |
| R11           | 110           | 600                 | 1100                 | 3                |
| R12           | 120           | 650                 | 1200                 | 3                |
| RESOURCE NAME | MIPS | BANDWIDTH | MEMORY SIZE | GRANULARITY |
|---------------|------|-----------|-------------|-------------|
| R13           | 130  | 700       | 1300        | 3           |
| R14           | 140  | 750       | 1400        | 3           |
| R15           | 150  | 800       | 1500        | 3           |
| R16           | 160  | 850       | 1600        | 3           |

Table registered resources (R12 to R16)

Comparison between the algorithms according to their processing time

The proposed secure resource and Job grouping based scheduling provides significant improvement in decreasing the processing time of jobs compared to dynamic job grouping strategy.

| NO of Jobs | processing time |
|------------|-----------------|
| SRJMB      | DJGB            |
| 3          | 4               | 10           |
| 5          | 7               | 22           |
| 8          | 11              | 28           |
| 10         | 16              | 34           |
| 14         | 21              | 40           |

Table Comparison between the algorithms

Conclusion and future scope:

In order to achieve secure resource and job scheduling including minimum processing time and maximum resource utilization, A Secure Resource RSA algorithm and Job Scheduling model with Job Grouping strategy in Grid Computing has proposed.

The Secure Resource and job scheduling algorithm will show better comparative result than dynamic job grouping based scheduling algorithm. The simulation environment has shown the proposed model is able to perform job and resource scheduling in grid environment and provide real grid Computing environment.

In future this work can be extended to implement job resource scheduling model with load balancing to increase the performance of Smart grid system.

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