Data Article

Benchmark dataset for the convex hull of 2D disks

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A B S T R A C T

In this paper, we present a benchmark dataset which can be used to evaluate the algorithms to construct the convex hull of 2D disks. The dataset contains disk arrangements including general and extremely biased cases, which are generated by a C++ program. The dataset is related to an article: "QuickHullDisk: A Faster Convex Hull Algorithm for Disks" in which the QuickHullDisk algorithm is presented and compared to the incremental algorithm which was reported by Devillers and Golin in 1995 [1].

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Convex hull is one of the most fundamental constructs in geometry and many algorithms to construct the convex hull have been extensively studied. Here, we present a benchmark dataset for the convex hull of 2D disks, which contains four types of disk arrangements spanning from general to extremely biased cases.

1.1. Four types of disk arrangements

1) RANDOM

A set \( \{D_{ij}\} \) where \( D_{ij} \) contains \( 10,000 \times i \) random disks and represents a data file. Thus, there are ten different files with the same number of disks. We randomly placed the disks of \( D_{ij} \) within a circular container. The container center is at the origin which has a sufficiently large radius so that the packing ratio \( \rho \) is maintained approximately 0.1 unless otherwise stated. Note that \( \rho \) is the ratio of the union of the area of individual disks to the area of the container. In RANDOM (and the other two test sets ON-BNDRY and MIXED), each disk has a random radius \( r \in [1.0, 10.0] \), two disks may intersect each other, and one may include another.
2) **ON-BNDRY**

A set \( ON-BNDRY = \{D_i|i = 1, 2, \ldots, 10\} \) where \( D_i \) contains \( 10,000 \times i \) disks touching a circular container from inside.

3) **MIXED**

A set generated by hybrid of RANDOM and ON-BNDRY: \( MIXED = \{D_{ij}|i, j = 1, 2, \ldots, 10\} \) where \( D_{ij} \) contains \( 1000 \times i \times j \) disks (out of \( 10,000 \times j \) disks) touching a circular container from inside. The remaining \( 1000 \times (10 - i) \times j \) disks not touching the container are randomly positioned within the container. Hence, a container touching ratio \( \gamma \) is well defined as the ratio of the number of disks touching the container to the total number of the disks.

4) **ON-A-LINE**

A set of congruent disks centered on a line (known as a sausage configuration [2]): \( ON-A-LINE = \{D_i|i = 1, 2, \ldots, 20\} \) where \( D_i \) contains \( 1000 \times i \) disks on the linear grid from left to right. Each disk \( d \in D_i \) has a unit radius (1.0) and the distance between the boundaries of two neighbour disks is kept 0.5.

Fig. 1 [1] shows examples of the four types of disk arrangements with reduced number of disks.

### 1.2. File format

For RANDOM, ON-BNDRY, ON-A-LINE, same file name convention is applied. Each file is named as ‘\( Nx.txt \)’ where ‘\( x \)’ denotes the total number of disks in the file. For example, a file ‘\( N10000.txt \)’ of RANDOM has 10,000 random disks. The file name of MIXED type includes the percentage of disks touching a container as well as total number of disks. A file ‘\( N10000_10.txt \)’ in MIXED includes 10% of the total 10,000 disks as disks touching a container (1000 disks).

The file format is as shown in Table 1 where we added descriptions of row and column. Fig. 2 shows an example of a data file which corresponds to Table 1. The first line contains an integer denoting the number of disks in the file. Then, each following line contains the definition of each disk with 4 fields: disk id (integer); x- and y-coordinate of disk center (floating-point number); radius of disk (floating-point number). The floating-point number is represented up to six decimal places.

![Fig. 1. Four types of disk arrangements (30 disks for each) (a) RANDOM (b) ON-BNDRY (c) MIXED (d) ON-A-LINE.](image)
For RANDOM dataset, disks are stored in the generation order. For ON-BNDRY dataset, the rightmost disk touching a circular container appears first in a file and the other disks in counterclockwise (CCW) orientation follow. For MIXED dataset, the disks touching a container appears in a file like ON-BNDRY and then random disks follow. For ON-A-LINE dataset, the leftmost disk appears first and the others follow from left to right.

2. Experimental design, materials, and methods

We generate a benchmark dataset by a code written by C++ language in Window 10 operating system. The code’s input parameters are data type (T), the number of disks (N), minimum radius (\(r_{\text{min}}\)), maximum radius (\(r_{\text{max}}\)) of disks, packing ratio (\(\rho\)), container touching ratio (\(\gamma\)) and distance between the boundaries of two neighbour disks (\(\delta\)).

### Table 1
An example for file format when ten disks are stored.

| Number of disks: 10 |
|---------------------|
| Id | X-coord. of center | Y-coord. of center | Radius |
|----|------------------|------------------|--------|
| 1  | 0.000000         | 0.000000         | 1.000000 |
| 2  | 2.500000         | 0.000000         | 1.000000 |
| 3  | 5.000000         | 0.000000         | 1.000000 |
| 4  | 7.500000         | 0.000000         | 1.000000 |
| 5  | 10.000000        | 0.000000         | 1.000000 |
| 6  | 12.500000        | 0.000000         | 1.000000 |
| 7  | 15.000000        | 0.000000         | 1.000000 |
| 8  | 17.500000        | 0.000000         | 1.000000 |
| 9  | 20.000000        | 0.000000         | 1.000000 |
| 10 | 22.500000        | 0.000000         | 1.000000 |

Fig. 2. A data file corresponding to Table 1.

For RANDOM dataset, disks are stored in the generation order. For ON-BNDRY dataset, the rightmost disk touching a circular container appears first in a file and the other disks in counterclockwise (CCW) orientation follow. For MIXED dataset, the disks touching a container appears in a file like ON-BNDRY and then random disks follow. For ON-A-LINE dataset, the leftmost disk appears first and the others follow from left to right.

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2.1. Disks in circular container

We fixed \(\rho\), \(r_{\text{min}}\) and \(r_{\text{max}}\) as 0.1, 1.0 and 10.0, respectively. The code starts from generating a circular container C centered at the origin, whose radius \(R\) is calculated by the following (1) and (2).

\[
S = \left\{ r_i \mid r_{\text{min}} + (r_{\text{max}} - r_{\text{min}}) \ast \left(\frac{i}{N + 1}\right), i = 1, 2, \ldots, N \right\}
\]  
(1)
\[
\rho = \frac{\cup (\text{Area of Disks})}{\text{(Container Area)}} = \frac{\sum \pi r_i^2}{\pi R^2}
\]  
(2)

Note. The generated set \( S \) is used only for calculating the radius of container but not for random disks.

1) RANDOM

Each disk is randomly generated in an axis-aligned bounding box of the container \( C \) (i.e. random location in the box and random radius between 1.0 and 10.0). If a generated disk is completely in the circular container \( C \), it is stored in a list. If not, a new disk with same radius is generated and tested again. Note that random number generator is used for generating random location and radius with current time as a seed unless otherwise stated.

2) ON-BNDRY

Each disk having random radius is generated to touch the container \( C \) from inside. The contact points to the container \( C \) are regularly spaced. The first generated disk is the rightmost extreme disk and the others follow in CCW order.

3) MIXED

First, disks touching container boundary \( \partial C \) are generated like ON-BNDRY. Then, disks not touching \( \partial C \) are generated like RANDOM. A random disk could touch \( \partial C \). To prevent this case, we choose a random disk within a shrunken container where the shrunken factor is 0.99.

Note. We confirmed that difference between computed packing ratio \( \rho \) of each data file and the fixed value (0.1) is less than \( 10^{-3} \). The contact points of disks to a container are different in \( 10^{-6} \) precision.

2.2. Disks in sausage configuration: ON-A-LINE

We set both \( r_{\text{min}} \) and \( r_{\text{max}} \) as 1.0. The first disk is generated to be centered at the origin and the other disks are generated to be centered on the \( x \)-axis where \( \delta \) is kept 0.5.

Note that all the data files is freely available from our dataset repository [3] which can be used by QuickhullDisk programs in the same repository and Voronoi Diagram Research Center, Hanyang University (http://voronoi.hanyang.ac.kr/).

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Conflict of Interest

The authors declare that they have no known competing financial interests or personal relationships that could have appeared to influence the work reported in this paper.

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