Flexibility of planar graphs without $C_4$ and $C_5^*$

Donglei Yang, Fan Yang†
School of Mathematics, Shandong University, Jinan 250100, China

Abstract

Let $G$ be a $\{C_4, C_5\}$-free planar graph with a list assignment $L$. Suppose a preferred color is given for some of the vertices. We prove that if all lists have size at least four, then there exists an $L$-coloring respecting at least a constant fraction of the preferences.

Key words: Planar graph, List coloring, Flexibility.

1 Introduction

In what follows, all graphs considered are simple, finite, and undirected, and we follow [II] for the terminologies and notation not defined here. A plane graph is a particular drawing of a planar graph in the Euclidean plane. Given a plane graph $G$, we denote its vertex set, edge set, face set and minimum degree by $V(G)$, $E(G)$, $F(G)$ and $\delta(G)$, respectively. The degree $d(v)$ is the number of edges incident with $v$. A vertex $v$ is called a $k$-vertex ($k^+$-vertex, or $k^-$-vertex) if $d(v) = k$ ($d(v) \geq k$, or $d(v) \leq k$, resp.). For any face $f \in F(G)$, the degree of $f$, denoted by $d(f)$, is the length of the shortest boundary walk of $f$, where each cut edge is counted twice. Analogously, a $k$-face ($k^+$-face, or $k^-$-face) is a face of degree $k$ (at least $k$, or at most $k$, resp.). We write $f = u_1u_2\ldots u_nu_1$ if $u_1, u_2, \ldots, u_n$ are the boundary vertices of $f$ in the clockwise order, and for integers $d_1, \ldots, d_n$, we say that $f$ is a $(d_1, \ldots, d_n)$-face if $d(u_i) = d_i$ for all $i \in \{1, 2, \ldots, n\}$. We say that $f$ is a $(d^+_1, \ldots, d^+_n)$-face if $d(v_1) \geq d_1$ and $d(v_i) = d_i$ for all $i \in \{2, \ldots, n\}$; and similarly for other combinations. Let $f_k(v)$, $n_k(v)$ and $n_k(f)$ denote the number of $k$-faces incident with the vertex $v$, the number of $k$-vertices adjacent to the vertex $v$, and the number of $k$-vertices incident with the face $f$, respectively. Moreover, we use $\delta(f)$ to refer to the minimum degree of vertices incident with $f$.

A list assignment $L$ for a graph $G$ is a function that to each vertex $v \in V(G)$ assigns a set $L(v)$ of colors, and an $L$-coloring is a proper coloring $\phi$ such that $\phi(v) \in L(v)$ for

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†Corresponding author. E-mail address: yangfan5262@163.com.
all $v \in V(G)$. If $G$ has an $L$-coloring, then we say that $G$ is $L$-colorable. Initiated by Dvořák, Norin, Postle [4], a request for a graph $G$ with a list assignment $L$ is a function $r$ with $\text{dom}(r) \subseteq V(G)$ such that $r(v) \in L(v)$ for all $v \in \text{dom}(r)$. For $\varepsilon > 0$, a request $r$ is $\varepsilon$-satisfiable if there exists an $L$-coloring $\phi$ of $G$ such that $\phi(v) = r(v)$ for at least $\varepsilon|\text{dom}(r)|$ vertices $v \in \text{dom}(r)$. We say that a graph $G$ with the list assignment $L$ is $\varepsilon$-flexible if every request is $\varepsilon$-satisfiable. Furthermore, we emphasize a stronger weighted form. A weighted request is a function $w$ that to each pair $(v, c)$ with $v \in V(G)$ and $c \in L(v)$ assigns a nonnegative real number. Let $w(G, L) = \sum_{v \in V(G), c \in L(v)} w(v, c)$. For $\varepsilon > 0$, we say that $w$ is $\varepsilon$-satisfiable if there exists an $L$-coloring $\phi$ of $G$ such that

$$\sum_{v \in V(G)} w(v, \phi(v)) \geq \varepsilon w(G, L).$$

We say that $G$ with the list assignment $L$ is weighted $\varepsilon$-flexible if every weighted request is $\varepsilon$-satisfiable.

It is worth pointing out that a request $r$ is 1-satisfiable if and only if the precoloring given by $r$ can be extended to an $L$-coloring of $G$. It is easy to deduce that weighted $\varepsilon$-flexibility implies $\varepsilon$-flexibility when we set there is only at most one color is requested at each vertex of $G$ and all such colors have the same weight 1.

Dvořák, Norin and Postle [4] proposed this topic and studied some natural questions in the context. They obtained some elementary results as follows, we say a list assignment $L$ is an $f$-assignment if $|L(v)| \geq f(v)$ for all $v \in V(H)$.

- There exists $\varepsilon > 0$ such that every planar graph with a 6-assignment is $\varepsilon$-flexible.
- There exists $\varepsilon > 0$ such that every planar graph of girth at least five with a 4-assignment is $\varepsilon$-flexible.
- For every integer $d \geq 2$, there exists $\varepsilon > 0$ such that every graph of maximum average degree at most $d$ and choosability at most $d - 1$ with a $(d + 2)$-assignment is weighted $\varepsilon$-flexible.

Especially, in [2], Dvořák et al. studied the 4-weighted flexibility of triangle-free planar graphs, and the result coincides with the choosability of triangle-free planar graphs. Later on, an interesting question is raised by Masařík [6] in the following.

**Question 1.** Does there exist $\varepsilon > 0$ such that every $C_4$-free planar graph $G$ with a 4-assignment is weighted $\varepsilon$-flexible?
If it is true, this would be optimal in terms of choosability \[5\]. However, it might be difficult to obtain such a result since even the procedure for getting the corresponding result of triangle-free is very involved. So far, triangle-free planar graphs are the only known result for \(\varepsilon\)-weighted flexibility with 4-assignment.

Based on the above results, we proceed a step towards Question 1 by proving the following theorem.

**Theorem 1.** There exists \(\varepsilon > 0\) such that every \(\{C_4, C_5\}\)-free planar graph with a 4-assignment is weighted \(\varepsilon\)-flexible.

In 2007, Voigt \[9\] proved that there exists \(\{C_4, C_5\}\)-free planar graphs which are not 3-choosable. Hence, our result is the best possible up to the list size.

The rest of the paper is organized as follows. In order to prove Theorem 1 in Section 2, we introduce some basic notation and essential tools used in list coloring settings. Using discharging method, we first find some necessary reducible configurations in Section 3, and then present our discharging rules and final analysis in Section 4.

### 2 Preliminaries

We use \(P_k\) to denote a path of order \(k\). Let \(H\) be a graph, \(S \subseteq V(H)\), \(S\) is called a \((P_3 + P_4)\)-independent set if \(H\) contains neither \(P_3\) nor \(P_4\) connecting two vertices in \(S\). Let \(1_S\) denote the characteristic function of \(S\), i.e., \(1_S(v) = 1\) if \(v \in S\) and \(1_S(v) = 0\) otherwise. For functions that assign integers to vertices of \(H\), we define addition and subtraction in the natural way, adding/subtracting their values at each vertex independently. For a function \(f : V(H) \to \mathbb{Z}\) and a vertex \(v \in V(H)\), let \(f \downarrow v\) denote the function such that \((f \downarrow v)(w) = f(w)\) for \(w \neq v\) and \((f \downarrow v)v = 1\).

Let \(G\) be a graph and \(H\) be an induced subgraph of another graph \(G\). For an integer \(k \geq 3\), let \(\delta_{G,k} : V(H) \to \mathbb{Z}\) be defined by \(\delta_{G,k}(v) = k - \deg_G(v)\) for each \(v \in V(H)\). Then \(H\) is said to be a \((P_3 + P_4,k)\)-reducible induced subgraph of \(G\) if

- **(FIX)** for every \(v \in V(H)\), \(H\) is \(L\)-colorable for every \(((\deg_H + \delta_{G,k}) \downarrow v)\)-assignment \(L\), and
- **(FORB)** for every \((P_3 + P_4, k)\)-independent set \(S\) in \(H\) of size at most \(k - 2\), \(H\) is \(L\)-colorable for every \(((\deg_H + \delta_{G,k} - 1_S)\)\)-assignment \(L\).

Note that **(FORB)** implies that \(\deg_H(v) + \delta_{G,k}(v) \geq 2\) for all \(v \in V(H)\).

To prove weighted \(\varepsilon\)-flexibility, we use the following observation made by Dvořák et al.\[4\]
Lemma 2.1 ([4]). Let $G$ be a graph and let $L$ be a list assignment for $G$. Suppose $G$ is $L$-colorable and there exists a probability distribution on $L$-colorings $\phi$ of $G$ such that for every $v \in V(G)$ and $c \in L(v)$, $\Pr[\phi(v) = c] \geq \varepsilon$. Then $G$ with $L$ is weighted $\varepsilon$-flexible.

Moreover, we use the following well-known result due to Thomassen [7].

Lemma 2.2 ([7]). Let $G$ be a connected graph and $L$ a list assignment such that $|L(u)| \geq \deg(u)$ for all $u \in V(G)$. If either there exists a vertex $u \in V(G)$ such that $|L(u)| > \deg(u)$, or some 2-connected component of $G$ is neither complete nor an odd cycle, then $G$ is $L$-colorable.

The following lemma provides an essential technique to deal with the weighted flexibility of graphs.

Lemma 2.3. For all integers $k \geq 3$, $b \geq 1$, there exists a positive constant $\varepsilon$ such that the following holds. Let $G$ be a $\{C_4, C_5\}$-free graph. If for every $Z \subseteq V(G)$, the graph $G[Z]$ contains an induced $(P_3 + P_4)$-reducible subgraph with at most $b$ vertices, then $G$ with any assignment of list of size $k$ is weighted $\varepsilon$-flexible.

Proof. Before proving the Lemma, we give the following claim at first.

Claim 2.4. For every integer $k \geq 3$, there exist $\varepsilon$, $\delta > 0$ as follows. Let $G$ be a $\{C_4, C_5\}$-free graph and $L$ be an assignment of lists of size $k$ to vertices of $G$. Then there exists a probability distribution on $L$-colorings $\phi$ of $G$ such that

(i) for all $v \in V(G)$ and $c \in L(v)$, we have $\Pr[\phi(v) = c] \geq \varepsilon$, and

(ii) for any $P_3 + P_4$-independent subset $S$ with $|S| \leq k - 2$, $\Pr[\phi(v) \neq c \text{ for all } v \in S] \geq \delta^{|S|}$.

Proof. Let $\delta = (\frac{1}{k})^b$, $\varepsilon = (\frac{1}{k})^{b+k-2}$. Assume $G$ is a $\{C_4, C_5\}$-free graph, $Z \subseteq V(G)$, there exists $Y \subseteq Z$ of size at most $b$ such that $G[Y] \subseteq G[Z]$ and $G[Y]$ is $(P_3 + P_4, k)$-reducible. We prove the claim by induction on $|V(G)|$. A random $L$-coloring $\phi$ of $G$ is chosen as follows: we choose an $L$-coloring $\phi_1$ of $G - Y$ at random from the probability distribution obtained by the induction hypothesis. Let $L'$ be the list assignment for $G[Y]$ defined by

$$L'(v) = L(v) \setminus \{\phi_1(u) : uv \in E(G), u \notin Y\}$$

for all $v \in Y$. Note that $|L'(v)| \geq \deg_{G[Y]}(v) + \delta_{G,k}(v)$ for all $v \in Y$, and thus $G[Y]$ has an $L'$-coloring by (FORB) applied with $S = \emptyset$. We choose an $L'$-coloring $\phi_0$ uniformly at random among all $L'$-colorings of $G[Y]$, and let $\phi$ be the union of the colorings $\phi_1$ and $\phi_0$. Note that $|L'(v)| \geq \deg_{G[Y]}(v) + \delta_{G,k}(v)$ for all $v \in Y$. 

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Firstly, we prove that (ii) holds. Let $S$ be a $(P_3 + P_4)$-independent subset in $G$, $S = S_1 \cup S_2$, where $S_1 = S \cap (G - Y)$, $S_2 = S \cap Y$. Obviously, both $S_1$ and $S_2$ are $P_3 + P_4$-independent subsets in $G - Y$ and $Y$, respectively. By the induction hypothesis for $G - Y$, we obtain $\Pr[(\forall v \in S_1) \phi(v) \neq c] = \Pr[(\forall v \in S_1) \phi_1(v) \neq c] \geq \delta^{\lfloor |S_1| \rfloor}$. If $S_1 = S$, then (ii) holds. Therefore, suppose that $|S_1| \leq |S| - 1$. Fix $\phi_1$, then consider the probability that $\phi_0$ gives all vertices of $S_2$ colors different from $c$. For $v \in S_2$, let $L_c(v) = L'(v) \setminus \{c\}$, and for $v \in Y \setminus S_2$, let $L_c(v) = L'(v)$. Note that $|L_c(v)| \geq \deg_{G[Y]}(v) + \delta_{G,k}(v) - 1_{S}(v)$ for all $v \in Y$. By (FORB), there exists an $L_c$-coloring of $G[Y]$ in that no vertex of $S_2$ is assigned color $c$. Since $\phi_0$ is chosen uniformly among the at most $k^b L'$-colorings of $G[Y]$, we conclude that the probability that no vertex of $S_2$ is assigned color $c$ by $\phi_0$ is at least $(\frac{1}{k})^b = \delta$. Consequently, under the assumption that $\phi_1$ does not assign color $c$ to any vertex of $S_1$, the probability that $\phi_0$ does not assign color $c$ to any vertex of $S_2$ is at least $\delta$. Hence,

$$\Pr[\phi(v) \neq c, \forall v \in S] = \Pr[\phi_1(u) \neq c, \forall u \in S_1; \phi_0(w) \neq c, \forall w \in S_2] \geq (\frac{1}{k})^b \delta^{|S_1|} \geq \delta^{|S_1| + 1} \geq \delta^{|S|}.$$ 

as required.

Next, we prove that (i) holds. Consider any vertex $v \in V(G)$ and a color $c \in L(v)$. If $v \in V(G - Y)$, then $\Pr[\phi(v) = c] = \Pr[\phi_1(v) = c] \geq \varepsilon$ by the induction hypothesis for $G - Y$. Therefore, we assume that $v \in Y$. Let $S$ be the set of neighbors of $v$ in $V(G) \setminus Y$. Since $G$ is $\{C_4, C_5\}$-free and all vertices in $S$ has a common neighbor, $S$ is $P_3 + P_4$-independent in $G - Y$. Furthermore, (FORB) implies $1 \leq \deg_{G[Y]}(v) + \delta_{G,k}(v) - 1_{v}(v) = \deg_{G[Y]}(v) + k - \deg_{G}(v) - 1 = k - 1 - |S|$, thus $|S| \leq k - 2$. By (ii) we obtain that no vertex of $S$ is assigned color $c$ by $\phi_1$ with probability at least $\delta^{k-2}$. (FIX) implies that there exists an $L'$-coloring of $G[Y]$ in which $v$ is assigned color $c$. Since $\phi_0$ is chosen uniformly among all $L'$-colorings of $G[Y]$, it follows that under the assumption that no vertex of $S$ is colored by $c$, the probability that $\phi(v) = c$ is at least $\delta$. Therefore,

$$\Pr[\phi(v) = c \forall v \in V(G)] = \Pr[\phi_1(w) \neq c \text{ for } \forall w \in S; \phi_0(u) = c \forall u \in Y] \geq (\frac{1}{k})^b \delta^{|S|} \geq (\frac{1}{k})^{b+k-2}.$$ 

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This completes the proof of Claim 2.4. □

Combining Claim 2.4 with Lemma 2.1, we obtain that Lemma 2.3 holds. □

Recall that a block of $H$ is a maximal connected subgraph without a cutvertex. The following lemma is easy to obtain by induction on the number of blocks in $H$. So we omit the proof.

**Lemma 2.5.** Let $L$ be an $f$-assignment for $H$. If every block $B$ of $H$ is $L'$-colorable for any $f \downarrow v$-assignment $L'$ with any fixed vertex $v \in V(B)$, then $H$ satisfies (FIX).

3 Reducible subgraphs

Before proceeding, we introduce the following notation. A vertex $v$ with $4 \leq d(v) \leq 12$ is bad if (i) the number of $(3, 4^-, v)$-face is $\lfloor \frac{d(v)}{3} \rfloor$; (ii) $f_3(v) = \lfloor \frac{d(v)}{2} \rfloor$. A vertex $v$ with $d(v) = 4$ is vice if $f_3(v) = 2$. And a vertex $v$ with $5 \leq d(v) \leq 12$ is dangerous if $f_{3,3}(v) = \lfloor \frac{d(v)-3}{2} \rfloor$. On the other hand, each $(3, 3, 4)$-face is called poor, each $(3, 3, v)$-face is called worse and each $(3, 3, v)$-face is called worst. Furthermore, denote by $f_{3,3}(v)$ (or $f_{3,4}(v)$), $f_{3b}(v)$ (or $f_{4b}(v)$) the number of worst 3-face (or worse 3-face incident with $v$), the number of $(3, w, v)$-face (or $(4, w, v)$-face) in which $w$ is a bad vertex, respectively. In particular, let $f_{3b}(v)$ be the number of $(v, w_1, w_2)$-face in which both $w_1$ and $w_2$ are bad vertices. For each $6^+$-face $f$, denote by $\xi(f)$ the number of poor 3-face sharing an edge with $f$. A nice path connecting $v$ and some vertex $u \in V(f)$ is a path of length at most two such that either

(i) $d(u) = 3$ and all internal vertices have degree 3 in $G'$; or

(ii) $d(u) = 4$ and all internal vertices are vice 4-vertices, moreover, every two consecutive 4-vertices in the path are contained in a $(4, 4, 4)$-face.

Note that in all figures of the paper, any vertex marked with • has no edges of $G$ incident with it other than those shown, any vertex marked with ■ is bad unless stated otherwise.

Here, we describe a quite general case of $(P_3 + P_4, 4)$-reducible configurations. Let $G$ be a $\{C_4, C_5\}$-free planar graph and $v$ a vertex of $G$. A $v$-stalk is one of the following subgraphs:

(a) An edge $vu_1$;

(b) A path $vu_1u_2$;

(c) A cycle $vu_1u_2$;
(d) A cycle $vu_1v_1$;
(e) A path $vu_1u_2$ and a path $vv_1u_1$;
(f) A cycle $vu_1w_1$;
(g) A path $vu_1u_2$ and a path $vw_1u_1$;
(h) A cycle $vw_1w_2$;
(i) A cycle $vv_1u_1$;
(j) A path $vv_1v_2v_3$, a path $vw_1v_1$ and a path $v_2u_1v_3$ ($v_1, v_2$ may be the same vertex);
(k) A path $vv_1v_3v_4$, a path $vv_2v_1$ and a path $v_3u_1v_4$ ($v_1, v_3$ may be the same vertex);
(l) A path $vv_1v_3v_4$, a path $v_3u_1v_4$, a path $v_2v_5v_6$, a path $v_5u_2v_6$ and an edge $v_1v_2$ ($v_1, v_3$ may be the same vertex, $v_2, v_5$ may be the same vertex);
(m) A path $vv_1v_2v_3$ and a path $v_2u_1v_3$ ($v_1, v_2$ may be the same vertex),

where $d(u_i) = 3$, $d(v_i) = 4$, $w_i$ is bad and $5 \leq d(w_i) \leq 12$ for each $i \in \{1, 2, \ldots, 6\}$. It is worth noting that for each bad vertex $w_i$, the $v$-stalk includes all neighbors of $w_i$ lying on the worse or worst 3-faces.

Lemma 3.1. Assume that $G$ is a $\{C_4, C_5\}$-free planar graph and $v \in V(G)$ with $5 \leq d(v) \leq 12$. Let $A = \{(a), (b), (c), (d), (e), (f), (g), (h), (j), (l)\}$, $B = \{(a), (b), (c), (d), (e), (l)\}$, $C = \{(b), (c), (e), (l)\}$, $D = \{(a), (b), (c), (e), (f), (g), (h), (j), (l)\}$ and $K = \{(i), (k)\}$ be the sets of stalks of $v$. If there exists a induced subgraph $H$ containing $v$ such that one of the following holds:

1. $V(H)$ either consists of the vertices lying on any combination of elements in $A$, or any combination of elements in $B$ together with one copy of (i), or any combination of elements in $C$ together with one copy of (k), or any combination of elements in $D$ together with one copy of (m) such that the resulting $(\deg_H + \delta_{G,A})$-assignment $L$ satisfies $|L(v)| \geq 3$.

2. $V(H)$ consists of the vertices lying on any combination of elements in $A$ together with only one element in $K$, except the special case where (d) and (k) appear in the same combination, such that the resulting $(\deg_H + \delta_{G,A})$-assignment $L$ satisfies $|L(v)| \geq 4$.

Then $G$ contains a $(P_3 + P_4, 4)$-reducible induced subgraph with at most 138 vertices.
Proof. For $x \in N(v)$, let $T_x$ be a $v$-stalk witnessing this case. Let $H$ be the subgraph of $G$ induced by $\bigcup_{x \in N(v)} T_x$. Clearly, $|V(H)| \leq (2d(w) - 1) \times \left\lfloor \frac{d(w)}{2} \right\rfloor \leq 138$. Thus it suffices to show that $H$ is $(P_3 + P_4, 4)$-reducible.

By $\{C_4, C_5\}$-freeness, it is easy to obtain that all stalks are pairwise vertex disjoint and there are no edges among them. Consider any vertex $z \in V(H)$ and a $(\deg_H + \delta_{G,A}) \downarrow z$-assignment $L$ for $H$, $H$ is $L$-colorable by Lemma 2.5 implying (FIX). Thus it only remains to show that $H$ satisfies (FORB). Let $S = \{s_1, s_2\}$ be a $(P_3 + P_4)$-independent subset in $G$ and $L'$ a $(\deg_H + \delta - 1S)$-assignment for $H$. First, we performing the following operations:

- Arbitrarily choose any two $v$-stalks $T_1$ and $T_2$ which containing at least a bad vertex or a nice path randomly;
- Arbitrarily choose a vertex $s_i$ in each $T_i$ for $i \in \{1, 2\}$ with $\text{dist}(s_1, s_2) = 1$ or $\text{dist}(s_1, s_2) \geq 4$.

When $|L(v)| \geq 3$, if $H$ is induced by the vertices of any combination of elements in $A$, let $A_0 = \{(a), (b), (c), (d)\}$. Obviously, the result holds for $A_0$. Next, we add (e), (f), (g), (h), (j), (l) into $A_0$ in turn to obtain $A$ and prove that the result also admits for $A$. Let $A_1 = A_0 \cup \{(e)\}$, we only need to consider whether (FORB) holds for the combination (b) and (e). In this situation, we can first $L'$-color $s_1, s_2, v$ in order and then greedily $L'$-color the remaining vertices. Hence, $A_1$ admits the result. Let $A_2 = A_1 \cup \{(f)\}$, we perform the above operation to define $S$, after that we $L'$-color the blocks which containing $s_1, s_2, v$ respectively, this is possible since $|L'| \geq 3$, and then greedily $L'$-color the remaining vertices. By the same argument, we eventually find that $A$ admits the result.

If $H$ is induced by the vertices of any combination of elements in $B$ as well as only one copy of (i), then we first greedily $L'$-color the vertices around the bad vertex $w_1$ in (i) and then greedily $L'$-color the remaining vertices.

Similarly, if $H$ is induced by the vertices of any combination of elements in $C$ as well as only one copy of (k) or by the vertices of any combinations of elements in $D$ as well as only one copy of (m), it is easy to verify the result also holds.

When $|L(v)| \geq 4$ and $H$ is induced by the vertices of any combination of elements in $A$ together with only one element in $K$. As usual, we perform the above operations. Then we first $L'$-color the vertices in the block which containing $s_i$, then greedily $L'$-color the remaining vertices. It is possible since $|L'(v)| \geq 4$. Hence, in all the cases, (FORB) holds.

In conclusion, $H$ is $(P_3 + P_4, 4)$-reducible. \[\square\]

Lemma 3.2. Let $G$ be a $\{C_4, C_5\}$-free planar graph. If $G$ contains one of the following configurations (see Figure 1). Here, $d(u_i) = 3$, $d(v_i) = 4$ and $5 \leq d(w_i) \leq 12$ for all $i$. 

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(1) A 4-vertex $v$ with $f_{bb}(v) = 1$;

(2) A 6-vertex $v$ with $f_{3,3}(v) = 1$, $f_{3b} = 1$;

(3) A $2t$-vertex $v$ with $f_{3,3}(v) = t - 1$, where $2 \leq t \leq 6$;

(4) A $2t$-vertex $v$ with $f_{3,3}(v) = t - 2$, $f_{bb}(v) = 1$, where $2 \leq t \leq 6$;

(5) A 5-vertex $v$ with $f_{3b} = 1$, and a path $vu_2u_3$;

(6) A 6-vertex $v$ with $f_{3b} = 1$, and two paths $vu_2u_3$, $vu_4u_5$.

Then $G$ contains a $(P_3 + P_4, 4)$-reducible induced subgraph with at most 29 vertices.

\[\text{Figure 1: Situation in Lemma 3.2}\]

**Proof.** Let $H$ be the subgraph of $G$ induced by $\{v, u_i, v_i, w_i\}$ for all $i$ and all neighbors of bad vertices lying on any the worst or worse $3$-faces. Note that $|V(H)| \leq 2d(w) - 1 + 6 \leq 29$. Consider any vertex $z \in V(H)$ and a $(\deg_H + \delta_G) \downarrow z$-assignment $L$ for $H$. Since each block $B$ in $H$ is a $3$-face or an edge and then we can easily observe that $B$ satisfies (\textbf{FIX}) under $L$, by Lemma 2.5, we know that $H$ satisfies (\textbf{FIX}).

Let $S$ be a $(P_3 + P_4)$-independent subset in $G$ with $|S| \leq 2$ and $L'$ a $(\deg_H + \delta - 1) S$-assignment for $H$. In (1), (2), (3), (4), since every two vertices is connected by a path of length 2 or 3 in $H$, (\textbf{FORB}) is implied by (\textbf{FIX}). In (5), if $S = \{v, u_2\}$ (or $S = \{u_2, u_3\}$), then we can first $L'$-color $v, u_2, u_3$ (or $u_3, u_2, v$) and then greedily $L'$-color the remaining vertices. If $S = \{s_1, u_3\}$, where $s_1$ lies on the worst or worse $3$-face $w_1s_1z$ incident with $w_1$, then we can first $L'$-color $u_3, s_1, z, v, u_2$ in order and the greedily $L'$-color the remaining vertices. In (6), the result also holds by the same argument.

Hence, both (\textbf{FIX}) and (\textbf{FORB}) hold, and thus $H$ is $(P_3 + P_4, 4)$-reducible. \hfill $\Box$
Lemma 3.3. Let $G$ be a $\{C_4, C_5\}$-free planar graph, there are three 3-faces $uvw, v_1v_2v_3, \ w_1w_2w_3$ that are incident with a 6-face $f = vww_1w_3v_3v_1v$ such that $d(u) = 3$, $d(v) = d(w) = d(v_1) = d(w_1) = 4$, and $d(v_i) \leq 12, d(w_i) \leq 12$ for all $i \in \{2, 3\}$. If each $x \in \{v_2, v_3, w_2, w_3\}$ satisfies that either $3 \leq d(x) \leq 4$ or $x$ is dangerous, then $G$ contains a $(P_3 + P_4, 4)$-reducible induced subgraph with at most 49 vertices.

![Figure 2: Situation in Lemma 3.3 and Lemma 3.4](image)

Proof. Let $H$ be the subgraph of $G$ induced by $A = \{u, v, w, v_1, v_2, v_3, w_1, w_2, w_3\}$ and all possible 3-neighbors inside a worst triangle incident with each $x \in \{v_2, v_3, w_2, w_3\}$, which is denoted by $B$ (possibly empty). Clearly, $V(H) = A \cup B$ and $|V(H)| \leq 49$. Now we are ready to prove that $H$ is $(P_3 + P_4, 4)$-reducible. It suffices to verify that $H$ satisfies (FIX) and (FORB). Firstly, consider an arbitrary $(\deg_H + \delta_{G,4})$-assignment $L$ for $H$. Then $|L(x)| > d_H(x)$ for each $x \in B \cup \{u\}$. In particular, $|L(x)| \geq d_{H'}(x)$ for each $x \in A$ and $|L(u)| > d_{H'}(u)$, where we denote $H' = H[A]$. It follows that (FIX) is immediately implied by Lemma 2.5.

It remains to verify that for all $\{P_3 + P_4\}$-independent set $S$, $H$ is $L'$-colorable for any $(\deg_H + \delta - 1_S)$-assignment $L'$. By the assumption, we can easily observe that either $|S \cap A| \leq 1$ or $S \in \{\{v, v_1\}, \{w, w_1\}, \{v_3, v_3\}\}$. The case when $S$ is a singleton is easily implied by (FIX). If $|S| = 2$ and $S$ contains at least one 3-vertex inside $B$, then we can first $L'$-color $H'$ by (FIX) and then greedily $L'$-color all vertices in $B$. Hence it remains to consider the case when $S \in \{\{v, v_1\}, \{w, w_1\}, \{v_3, v_3\}\}$. In all cases, we firstly $L'$-color $H[S]$ and let $\{a, b\}$ be the resulting colors on $S$. Let $H^* = H' - S$ and define a list assignment $L^*$ such that $L^*(z) = L'(z) \setminus \{a, b\}$ for each $z \in V(H') \cap N_H(S)$, while $L^*(z) = L'(z)$ for each $z \in V(H') \setminus N_H(S)$. It is easy to see that $|L^*(z)| \geq \deg_{H^*}(z)$ for all $z \in V(H^*)$. Since $|L^*(u)| > \deg_{H^*}(u)$, by Lemma 2.2 we can always $L^*$-color all vertices in $A - S$ and then extend the coloring to $B$ greedily. Hence $H$ satisfies (FORB).

This completes the proof of Lemma 3.3.

□
Lemma 3.4. Let \( G \) be a \( \{C_4, C_5\} \)-free planar graph, if there are two 3-faces \( uvw, x_1x_2x_3 \) that are adjacent to a 6-face \( uwx_1x_3x_4x_5u \) such that \( d(u) = d(x_5) = 3 \) and \( d(v) = d(w) = d(x_1) = 4 \). If each \( x \in \{x_2, x_3, x_4\} \) satisfies that either \( 3 \leq d(x) \leq 4 \) or \( x \) is dangerous, then \( G \) contains a \( (P_3 + P_4, 4) \)-reducible induced subgraph with at most 37 vertices.

Proof. Let \( H \) be the subgraph of \( G \) induced by \( A = \{u, v, w, x_1, x_2, x_3, x_4, x_5\} \) and all possible 3-neighbors inside a worst triangle incident with each \( x_i \) for all \( i \in \{2, 3, 4\} \), which is denoted by \( B \). Clearly, \( V(H) = A \cup B \) and \( |V(H)| \leq 37 \). Now we are ready to prove that \( H \) is \( (P_3 + P_4, 4) \)-reducible. It suffices to verify that \( H \) satisfies (\text{FIX}) and (\text{FORB}). Consider an arbitrary \( (\text{deg}_H + \delta_{C,4}) \)-assignment \( L \) of \( H \). If \( |L(x)| > d_H(x) \) for each \( x \in B \cup \{u, x_5\} \), then \( |L(x)| \geq d_H(x) \) for each \( x \in A \), where we denote \( H' = H[A] \). It follows that (\text{FIX}) is immediately implied by Lemma 2.5.

It remains to verify that for all \( \{P_3 + P_4\} \)-independent set \( S \), \( H \) is \( L' \)-colorable for any \( (\text{deg}_H + \delta - 1_S) \)-assignment \( L' \). By the assumption, we can easily observe that either \( |S \cap A| \leq 1 \) or \( S \in \{\{w, x_1\}, \{x_3, x_4\}, \{x_4, x_5\}, \{u, x_5\}\} \). The case when \( S \) is a singleton is easily implied by (\text{FIX}). If \( |S| = 2 \) and \( S \) contains at least one 3-vertex inside \( B \), then we can first \( L' \)-color \( H' \) by (\text{FIX}) and then greedily \( L' \)-color all vertices in \( B \). Hence it only remains to consider the case when \( S \in \{\{w, x_1\}, \{x_3, x_4\}, \{x_4, x_5\}, \{u, x_5\}\} \).

If \( S = \{u, x_5\} \), then consider the subgraph \( H' \) together with the list assignment \( L' \), by Lemma 2.2 we can always \( L' \)-color all vertices in \( A \), and then greedily \( L' \)-color all vertices in \( B \). In the remaining cases, we firstly \( L' \)-color \( H[S] \) and let \( \{a, b\} \) be the resulting colors on \( S \). Define a list assignment \( L^* \) on \( V(H) - S \) such that \( L^*(z) = L'(z) \setminus \{a, b\} \) for each \( z \in N_H(S) \), while \( L^*(z) = L'(z) \) for each \( z \in V(H) \setminus N_H(S) \). Since \( |L^*(u)| = 4, |L^*(x_5)| = 3 \), by Lemma 2.2 we can always \( L^* \)-color all vertices in \( A - S \) and then extend the coloring to \( B \) greedily. Hence \( H \) satisfies (\text{FORB}).

\[
\begin{array}{c}
\text{Figure 3: Situation in Lemma 3.5}
\end{array}
\]

Lemma 3.5. Let \( G \) be a \( \{C_4, C_5\} \)-free planar graph, if there are three 3-faces \( uvw, x_1x_2x_3, x_4x_5x_6 \) that are adjacent to a 6-face \( uwx_1x_3x_4x_5u \) such that \( d(u) = 3 \) and \( d(v) = d(w) = d(x_1) = 4 \).
If \( d(x_5) = 4 \) (see Figure 3). If \( x_5 \) is dangerous with \( d(x_5) \) odd, while each \( x \in \{x_2, x_3, x_4\} \) satisfies that either \( 3 \leq d(x) \leq 4 \) or \( x \) is dangerous, then \( G \) contains a \((P_3+P_4, 4)\)-reducible induced subgraph with at most 45 vertices.

**Proof.** Let \( H \) be the subgraph of \( G \) induced by \( A = \{u, v, w, x_1, x_2, x_3, x_4, x_5, x_6\} \) and all possible 3-neighbors inside a worst triangle incident with each \( x_i \) for all \( i \in \{2, 3, 4, 5\} \), which is denoted by \( B \). Clearly, \( V(H) = A \cup B \) and \( |V(H)| \leq 45 \). Now we are ready to prove that \( H \) is \((P_3 + P_4, 4)\)-reducible. First, consider a \((\deg H + \delta_G, 4)\)-assignment \( L \) for \( H \), it is easy to observe that \((\text{FIX})\) is implied by Lemma 2.5.

It remains to verify that for all \( \{P_3 + P_4\}\)-independent set \( S \), \( H \) is \( L' \)-colorable for any \((\deg H + \delta - 1_S)\)-assignment \( L' \). The case when \( S \) is a singleton is easily implied by \((\text{FIX})\), while the case when \( S \) contains a 3-vertex inside \( B \) is implied by \((\text{FIX})\) and Lemma 2.2. Hence it suffices to consider \( S \in \{\{w, x_1\}, \{x_3, x_4\}, \{u, x_5\}\} \).

If \( S = \{u, x_5\} \), then consider the subgraph \( H[A] \) together with the list assignment \( L' \), by Lemma 2.2 we can always \( L' \)-color all vertices in \( A \), and then greedily \( L' \)-color all vertices in \( B \). In the remaining cases, we firstly \( L' \)-color \( H[S] \) and let \( \{a, b\} \) be the resulting colors on \( S \). Let \( H^* = H - S \) and define a list assignment \( L^* \) such that \( L^*(z) = L'(z) \setminus \{a, b\} \) for each \( z \in V(H^*) \cap N_H(S) \), while \( L^*(z) = L'(z) \) for each \( z \in V(H^*) \setminus N_H(S) \). Since \( |L^*(u)| = |L^*(x_5)| = 4 \), by Lemma 2.2, we can always \( L^* \)-color all vertices in \( A - S \) and then extend the coloring to \( B \) greedily. Hence \( H \) satisfies \((\text{FORB})\).

![Figure 4: Situation in Lemma 3.6](image-url)
Lemma 3.6. If $G$ has a subgraph isomorphic to one of the configurations in Figure 4, then $G$ contains a $(P_3 + P_4, 4)$-reducible induced subgraph with at most 9 vertices.

Proof. Let $H$ be the graph isomorphic one of the configurations in Figure 4. Note that (FIX) is easily obtained by Lemma 2.5. It remains to verify that for all $(P_3 + P_4)$-independent set $S$, $H$ is $L'$-colorable for any $(\deg_H + \delta - 1_S)$-assignment $L'$. The case when $S$ is a singleton is easily implied by (FIX). Now we only consider $|S| = 2$. In all cases, $S$ induces a pendant edge in $H$ and it follows that we can first $L'$-color $H[S]$, and then greedily $L'$-color the remaining vertices in $H$. Hence, (FORB) holds.

![Figure 5: Situation in Lemma 3.7](image)

Lemma 3.7. Let $G$ be a planar graph satisfying any one of the following conditions (see Figure 5). Here $d(u_i) = 3$, $d(v_i) = 4$, $w_i$ is a bad $5^+$-vertex for all $i$, while $w$ is both bad and dangerous.

1. A 5-vertex $v$ with $f_3(v) = 2$, and $vu_2$ is a pendant edge, $vu_1x$ is a 3-face, where $z_1$ is either a 4-vertex or bad;

2. A 6-vertex $v$ with $f_3(v) = 3$, and $vu_1v_1$, $vu_2y$ are 3-faces, where $z_2$ is either a 4-vertex or bad;

3. A 6-vertex $v$ with $f_3(v) = 3$, and $f_{5b}(v) = 2$;

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(4) A 6-vertex $v$ with $f_3(v) = 3$, and $f_{3,3}(v) = 1$, $f_{4b}(v) = 1$;

(5) A 6-vertex $v$ with $f_3(v) = 3$, and $f_{4b}(v) = 1$, two paths $v_1v_2v_3$, $v_2u_3v_3$ ($v_1$, $v_2$ may be the same vertex);

(6) A 6-vertex $v$ with $f_3(v) = 3$, and $f_{3,3}(v) = 1$, four paths $vv_1v_2v_3$, $v_2u_3v_3$, $vv_4v_5v_6$, $v_5u_4v_6$ ($v_1$, $v_2$ may be the same vertex, $v_4$, $v_5$ may be the same vertex).

Then $G$ contains a $(P_3 + P_4, 4)$-reducible induced subgraph with at most 37 vertices.

Proof. We denote by $N^*(x)$ the set of neighbors of each bad vertex $x$ lying on the worse or worst 3-faces. In all cases, if $d(w)$ is even, then by the assumption that $f_{3,3}(w) = \frac{d(w)-2}{2}$, Lemma 3.2(3) directly implies a $(P_3 + P_4, 4)$-reducible induced subgraph on at most 37 vertices. Hence we may assume that $d(w)$ is odd.

Now consider the cases (1)-(4), let $H$ be the subgraph induced by $v, w$, all $u_i$, $v_i$, $w_i$, $z_i$ and all vertices in $N^*(x)$ for each bad vertex $x \in \{w, w_1, w_2, (z_i)\}$. Then we claim that $H$ is a $(P_3 + P_4, 4)$-reducible induced subgraph. In fact, for any $(\deg_H + \delta_{G,4})$-assignment $L$ of $H$, we have $|L(w)| = 2, |L(v)| \geq 3$ and $|L(u_i)| \geq 3, |L(v_i)| = 2$ for each $i \in [2]$. In addition, $|L(y)| = 3$ for all $y \in N^*(w), |L(u_2)| = 2$ in configuration (1), while $|L(u_i)| = 3$ for each $i \in [2]$ in configurations (2)-(4). It follows from Lemma 2.5 that (\textbf{FIX}) holds. It remains to verify that for all $\{P_3 + P_4\}$-independent set $S$, $H$ is $L'$-colorable for any $(\deg_H + \delta - 1_s)$-assignment $L'$. The case when $S$ is a singleton is easily implied by (\textbf{FIX}). Now we only consider $|S| = 2$. In all cases, we can first $L'$-color $H[S]$, and then greedily $L'$-color the remaining vertices in $H$. Hence, (\textbf{FORB}) holds.

Consider the configuration (5)-(6). Here, we just consider (6) since (5) can be solved by the same argument. Let $H$ be the subgraph induced by $v, w$, all $u_i, v_i$ and all vertices in $N^*(w)$. Now we claim that $H$ is a $(P_3 + P_4, 4)$-reducible induced subgraph. In fact, for any $(\deg_H + \delta_{G,4})$-assignment $L$ of $H$, we have $|L(x)| = 2$ for each $x \in \{w, v_3, v_6\}, |L(y)| = 3$ for each $y \in \{v, u_1, v_2, u_3, u_4, v_1, v_2, v_4, v_5\} \cup N^*(w)$. It follows from Lemma 2.5 that (\textbf{FIX}) holds. It remains to verify that for all $\{P_3 + P_4\}$-independent set $S$, $H$ is $L'$-colorable for any $(\deg_H + \delta - 1_s)$-assignment $L'$. By similar arguments, in all cases, we can first $L'$-color $H[S]$, and then greedily $L'$-color the remaining vertices in $H$. Hence, (\textbf{FORB}) holds.

\hfill \Box

\textbf{Lemma 3.8.} Let $uv$ be an edge of $G$ such that $d(u) = 3$ and $d(v) = 4$. If there exists a vice vertex $v_2$ such that both $vv_1v_2$ and $v_2v_3u_1$ are 3-faces incident with $v_2$ with $d(u_1) = 3$ and $d(v_i) = 4$ for each $i \in \{1, 2, 3\}$. Then $G$ contains a $(P_3 + P_4, 4)$-reducible induced subgraph with at most 6 vertices.
Proof. Let $H$ be the subgraph induced by $\{u, v, u_1, v_1, v_2, v_3\}$, consider any vertex $z \in V(H)$ and a $(\deg_H + \delta_{G,4}) \downarrow z$-assignment $L$ for $H$, $H - z$ is $L$-colorable by Lemma 2.2 implying (FIX). Now we only consider $|S| = 2$. Since $S$ induces a pendant edge in $H$ and it follows that we can first $L'$-color $H[S]$, and then greedily $L'$-color the remaining vertices in $H$. Consequently, (FORB) admits. Hence, $H$ is $(P_3 + P_4, 4)$- reducible. \hfill \square

4 Discharging Process

To prove Theorem 1 the main idea is to apply Lemma 2.3. Let $G$ be a $\{C_4, C_5\}$-free planar graph. If $G$ satisfies the assumption in Lemma 2.3, then we are done. Now we may assume that there exists $Z_0 \subseteq V(G)$ such that $G[Z_0]$ does not contain any $(P_3 + P_4, 4)$- reducible subgraph on at most 100 vertices. Let $G' = G[Z_0]$. Since $G'$ is also a plane graph, by Euler’s Formula, we obtain

$$\sum_{v \in Z_0} (d(v) - 2) + \sum_{f \in F(G')} (-2) = -4.$$ 

Now we redistribute the charges of all vertices and faces as follows, where we use $c(x \rightarrow y)$ to denote the charge sent from an element $x$ to another element $y$. Let $f$ be a face of $G'$ and $v_1, v_2, v_3$ be three vertices on $f$.

R0. If $f$ is a 3-face and $d(v_1) \geq 13$, then $c(v_1 \rightarrow f) = \frac{4}{3}$, $c(v_2 \rightarrow f) = c(v_3 \rightarrow f) = \frac{1}{3}$;

R1. For each $v \in V(f)$, if either $d(v) = 3$, or $d(v) \geq 4$ and $d(f) \geq 6$, then $c(v \rightarrow f) = \frac{1}{3}$;

R2. Let $f = v_1v_2v_3$ with $d(v_1) = 4$.

R2.1. If $d(v_2) = 4$, $d(v_3) = 3$ and $v_1$ is not vice, then $c(v_1 \rightarrow f) = 1$, $c(v_2 \rightarrow f) = \frac{2}{3}$;

R2.2. If $d(v_2) = 4$, $5 \leq d(v_3) \leq 12$ and $v_i$ is bad, $v_3-i$ is not vice for some $i \in \{1, 2\}$, then $c(v_i \rightarrow f) = \frac{1}{2}$, $c(v_{3-i} \rightarrow f) = \frac{5}{6}$;

R2.3. Otherwise, let $c(v_1 \rightarrow f) = \frac{2}{3}$.

R3. If $f = v_1v_2v_3$ with $3 \leq d(v_1) \leq d(v_2) \leq 4 < d(v_3) \leq 12$, then $c(v_3 \rightarrow f) = 2 - c(v_1 \rightarrow f) - c(v_2 \rightarrow f)$.

R4. Let $f = v_1v_2v_3$ such that $d(v_1) = 3$, $5 \leq d(v_2) \leq 12$ and $5 \leq d(v_3) \leq 12$.

R4.1. If there exists a bad vertex $v_i \in \{v_2, v_3\}$, then $c(v_i \rightarrow f) = \frac{2}{3}$, $c(v_{5-i} \rightarrow f) = 1$;

R4.2. If there is no bad vertices on $f$, then $c(v_2 \rightarrow f) = c(v_3 \rightarrow f) = \frac{5}{6}$.

R5. Let $f = v_1v_2v_3$ such that $d(v_1) = 4$, $5 \leq d(v_2) \leq 12$, and $5 \leq d(v_3) \leq 12$. 

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R5.1. If there exists a bad vertex $v_i \in \{v_2, v_3\}$, then $c(v_i \rightarrow f) = \frac{1}{2}$, $c(v_5-i \rightarrow f) = \frac{5}{6}$;

R5.2. If neither $v_2$ nor $v_3$ is bad, then $c(v_2 \rightarrow f) = c(v_3 \rightarrow f) = \frac{2}{3}$.

R6. If $f = v_1v_2v_3$ such that $5 \leq d(v_1) \leq d(v_2) \leq d(v_3) \leq 12$,

R6.1. If there exists exactly one vertex which is both bad and dangerous, say $v_1$, then $c(v_1 \rightarrow f) = \frac{1}{2}$, $c(v_2 \rightarrow f) = c(v_3 \rightarrow f) = \frac{3}{4}$;

R6.2. If there are two vertices, say $v_1, v_2$, where each $v_i$ is bad and dangerous, then $c(v_1 \rightarrow f) = c(v_2 \rightarrow f) = \frac{1}{2}$, $c(v_3 \rightarrow f) = 1$.

R6.3. Otherwise, let $c(v_i \rightarrow f) = \frac{2}{3}$ for each $i \in \{1, 2, 3\}$.

Let $ch(x)$ be the new charge for each $x \in V(G') \cup F(G')$ after applying rules R1-R6. It is easy to see that for each vertex $v \in V(G')$, $f_3(v) \leq \lfloor \frac{d(v)}{2} \rfloor$. Since $G'$ does not contain any $(P_3 + P_4, 4)$-reducible subgraph, Lemma 3.1 and Lemma 3.2 immediately imply the following corollary.

**Corollary 4.1.** Each $3^+$-vertex $v \in V(G')$ satisfies the following:

1. $f_{33}(v) + f_{34}(v) + f_{3b}(v) + f_{bb}(v) \leq \lfloor \frac{d(v)}{2} \rfloor - 1$;

2. If $v$ is dangerous, then $f_{3b}(v) = f_{bb}(v) = 0$. Furthermore, if $d(v)$ is odd, then $f_{34}(v) = f_{4b}(v) = 0$.

Moreover, we consider all $5^+$-vertices that satisfy certain properties as follows.

**Claim 4.2.** For each vertex $v$ of $G'$. If $d(v) \geq 5$ and $v$ is not dangerous, then $ch_1(v) \geq \frac{1}{6}$.

**Proof.** By Corollary 4.1(1), $f_{33}(v) + f_{34}(v) + f_{3b}(v) + f_{bb}(v) \leq \lfloor \frac{d(v)}{2} \rfloor - 1$. If $d(v) \geq 5$ and $v$ is not dangerous, then $f_{33}(v) \leq \max\{0, \frac{d(v)+1}{2} - 3\}$. Hence, when $d(v) = 5$, $ch_1(v) \geq 5 - 2 - 1 - \frac{5}{6} - \frac{3}{3} \geq \frac{1}{6}$. When $d(v) = 6$, $ch_1(v) \geq 6 - 2 - 1 - \frac{5}{6} - \frac{3}{3} \geq \frac{1}{6}$. When $d(v) \geq 7$, $ch_1(v) \geq d(v) - 2 - \frac{4}{3} \times (\frac{d(v)+1}{2} - 3) - 2 - \frac{1}{3} \times \frac{d(v)+1}{2}$

$$\geq d(v) - 5d(v) + 5\frac{6}{6}$$

$$= \frac{d(v) - 5}{6} \geq \frac{1}{5}.$$ 

$\Box$
Claim 4.3. For each $5^+$-vertex $v$ of $G'$. If $v$ is incident with a face $f = (u, v, w)$ in which $u$ is a bad $4$-vertex and $d(w) \geq 4$, then $c_1(v) \geq \frac{1}{6}$.

Proof. If $d(v) \geq 13$, then $c_1(v) \geq d(v) - 2 - \frac{4}{3} \times \left( \left\lfloor \frac{d(v)}{2} \right\rfloor - 1 \right) = \frac{2d(v) - \left\lfloor \frac{d(v)}{2} \right\rfloor - 2}{3} \geq \frac{d(v) - g}{6} \geq \frac{1}{6}$ by R1 and R6. Next it suffices to consider the case when $d(v) \leq 12$. Moreover, by Claim 4.2 it suffices to consider the case when $v$ is dangerous.

By Corollary 4.1(2), since $f_{3, 3}(v) = \frac{d(v) - 3}{2}$, so $f_{bb}(v) = 0$ and $w$ is not bad. Then $c(v \to f) \leq \frac{2}{3}$. If $d(v) = 5$, then $v$ is bad, and by Lemma 3.2(1), we know that $d(w) \geq 5$. Hence, by R5.1, $c_1(v) \geq 5 - 2 - 4 \frac{3}{3} \frac{1}{2} - \frac{2}{3} \frac{1}{3} \frac{1}{6}$. When $d(v) = 6$ and $f_{3, 3}(v) = 1$, if $v$ is bad, then $d(w) \geq 5$ and by R5.1, $c(v \to f) = \frac{1}{2}$ and it follows that $c_1(v) \geq 6 - 2 - 4 \frac{3}{3} \frac{1}{2} - \frac{1}{2} - \frac{3}{3} \frac{1}{6} \geq \frac{1}{6}$. Otherwise, since $f_{3, 4}(v) + f_{3b}(v) + f_{bb}(v) = 0, c_1(v) \geq 6 - 2 - 4 \frac{3}{3} \frac{5}{6} - \frac{2}{3} \frac{3}{3} \frac{1}{6} \geq \frac{1}{6}$. When $d(v) \geq 7$ is odd, by Corollary 4.1(2), we have $f_{3b}(v) + f_{3, 4}(v) + f_{bb}(v) = 0$ and

\[
\begin{align*}
\text{ch}_1(v) & \geq d(v) - 2 - \frac{4}{3} \times \left( \left\lfloor \frac{d(v)}{2} \right\rfloor - 2 \right) - \frac{2}{3} \times \frac{1}{3} \times \frac{d(v) + 1}{2} \\
& \geq d(v) - \frac{5d(v) + 5}{6} \\
& = \frac{d(v) - 5}{6} \\
& \geq \frac{1}{3}.
\end{align*}
\]

When $d(v) \geq 8$ is even, we have $f_{3, 3}(v) = \frac{d(v)}{2} - 2, f_{3, 4}(v) \leq 1$. Hence,

\[
\begin{align*}
\text{ch}_1(v) & \geq d(v) - 2 - \frac{4}{3} \times \left( \left\lfloor \frac{d(v)}{2} \right\rfloor - 2 \right) - 1 - \frac{2}{3} \times \frac{1}{3} \times \frac{d(v)}{2} \\
& \geq d(v) - \frac{5d(v)}{6} - 1 \\
& = \frac{d(v) - 6}{6} \\
& \geq \frac{1}{3}.
\end{align*}
\]

A vertex $v$ in $G'$ is called well when $\text{ch}_1(v) \geq \frac{1}{12}$. Given a poor face $f = (3, 4, 4)$ and a well vertex $v$. From now on, let $\varpi(v)$ be the number of nice paths starting at $v$. For each poor face $f$, we apply the following rules.

R7. If $f$ is poor and $g$ is a $7^+$-face sharing an edge with $f$, then $f$ receives $(\frac{d(g)}{3} - 2)/\xi(g)$ from $g$.

R8. If $f$ receives less than $\frac{1}{3}$ by R7 and $v_1, \ldots, v_t (1 \leq t \leq 2)$ are the well vertices such that all nice paths connecting each $v_i$ with $f$ has the same internal vertices, then $f$ receives $\frac{1}{62}$ from each $v_i (i \in \{1, \ldots, t\})$. 

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Let $ch_2(x)$ be the final charge for each element $x \in V(G') \cup F(G')$ after applying R7 and R8.

4.1 Each poor face $f$ satisfies $ch_2(f) \geq 0$.

Let $f = (u, v, w)$ be a poor face such that $d(u) = 3$, $d(v) = 4$, $d(w) = 4$ and $f_1, f_2, f_3$ be three adjacent faces sharing edge $vw, vu, uw$ with $f$, respectively. In addition, let $x_5$ be the neighbor of $u$ outside $f$. If $v$ is not vice, then by R1-R2, $ch_2(f) \geq 0$. Hence by symmetry, we may assume that both $v$ and $w$ are vice. Let $w_1, x_1, v_1$ be the other two neighbors of $w$ and $v$ outside $f$, respectively. Then they are all $4^+$-vertices.

If there exists either a $5^+$-vertex or a $4^+$-vertex that is not vice, among $\{v_1, y_1, x_1, w_1\}$, say $x_1$, then by Claim 4.3, $ch_1(x_1) \geq \frac{1}{6}$. In particular, if there are at least two such vertices as $x_1$ among $\{v_1, y_1, x_1, w_1\}$, then by R7-R8, we have $ch_2(f) \geq 0$. Hence, without loss of generality, we further assume that at least three among them, say $x_1, w_1, v_1$, are vice $4^+$-vertices. Note that the vertex sets $\{v_2, v_3, w_2, w_3\}$ and $\{x_1, x_2, x_3, x_4\}$ do not overlap.

Claim 4.4. If $d(f_1) = 6$ and $x_1, y_1, w_1, v_1$ are vice $4^+$-vertices, then $f$ receives at least $\frac{1}{6}$ from a vertex in $\{v_2, v_3, w_2, w_3\}$.

Claim 4.5. If $d(f_2) = 6$, then $f$ receives at least $\frac{1}{6}$ from $x_2, x_3, x_4, x_5$ altogether.
Proof. Suppose it is not true. We may similarly assume that each vertex in \{x_2, x_3, x_5\} has degree at most 12.

If \(d(x_5) = 3\) and \(d(x_4) \geq 13\), then by R8, \(f\) receives \(\frac{1}{6}\) from \(x_4\), a contradiction. If \(d(x_5) = 3\) and \(d(x_4) \leq 12\), then by Lemma 3.4 there exists a \(5^+\)-vertex in \(\{x_2, x_3, x_4\}\) that is not dangerous. It follows from Claim 4.2 and R8 that \(f\) receives \(\frac{1}{6}\) from \(x_2, x_3, x_4\) altogether.

If \(d(x_5) = 4\), then \(x_5\) is not vice. Moreover, by Corollary 4.1(1), \(f_{3,4}(v) = 0\), and it is easy to see that \(x_5\) is well. So \(f\) receives \(\frac{1}{6}\) from \(x_5\).

If \(5 \leq d(x_5) \leq 12\), then by Claim 4.2 it remains to consider the case when \(x_5\) is dangerous. Now we have the following observations:

1. \(f_{3b}(x_5) = f_{6b}(x_5) = 0\);
2. \(x_4x_5\) is contained in a 3-face, say \(x_4x_5x_6\);
3. \(x_4, x_6\) are \(4^+\)-vertices;
4. \(d(x_5)\) is odd.

In fact, (1) and (3) are immediately derived from Corollary 4.1. For (2), if \(x_4x_5\) is not contained in a 3-face, then \(x_5\) is incident with at most \(\frac{d(v)-2}{2}\) 3-faces. Since \(x_5\) is dangerous, it is easy to see that if \(d(x_5)\) is odd, then by Corollary 4.1(2), \(f_{3,3}(x_5) = 0\) and \(ch_1(x_5) \geq d(x_5) - 2 - \frac{4}{3} \times \frac{d(x_5)-3}{2} - \frac{1}{3} \times \lceil \frac{d(x_5)+1}{2} \rceil \geq \frac{1}{6}\), a contradiction. If \(d(x_5)\) is even, then \(ch_1(x_5) \geq d(x_5) - 2 - \frac{4}{3} \times \frac{d(x_5)-4}{2} - 1 - \frac{1}{3} \times \frac{d(x_5)+2}{2} \geq \frac{1}{3}\), a contradiction. For (4), If \(d(x_5)\) is even, then when \(wx_5\) is not contained in a 3-face, we know \(x_5\) is incident with at most \(\frac{d(v)-2}{2}\) 3-faces. Applying similar arguments as above, we know that \(ch_1(x_5) \geq d(x_5) - 2 - \frac{4}{3} \times \frac{d(x_5)-4}{2} - 1 - \frac{1}{3} \times \frac{d(x_5)+2}{2} \geq \frac{1}{3}\). In all those cases, \(f\) receives \(\frac{1}{6}\) from \(x_5\), a contradiction.

By the above observations, we know \(x_5\) is also a bad vertex. Consider the 3-face \(f' = x_4x_5x_6\), if \(d(x_i) \geq 5\) for all \(i \in \{4, 6\}\), then by R6, \(c(x_5 \rightarrow f') = \frac{1}{2}\) and it follows that \(ch_1(x_5) \geq d(x_5) - 2 - \frac{4}{3} \times \frac{d(x_5)-3}{2} - \frac{1}{2} - \frac{1}{3} \times \frac{d(x_5)+1}{2} \geq \frac{1}{6}\). If \(d(x_i) = 4\) for some \(i \in \{4, 6\}\), then \(x_{10-i}\) is not bad. If \(d(x_{10-i}) \geq 13\), then by R0, \(c(x_5 \rightarrow f') = \frac{1}{3}\). We know that \(ch_1(x_5) \geq d(x_5) - 2 - \frac{4}{3} \times \frac{d(x_5)-3}{2} - \frac{1}{2} - \frac{1}{3} \times \frac{d(x_5)+1}{2} \geq \frac{1}{6}\), a contradiction. If \(5 \leq d(x_{10-i}) \leq 12\), then by R5.1, \(c(x_5 \rightarrow f') = \frac{1}{2}\), and it follows that \(ch_1(x_5) \geq d(x_5) - 2 - \frac{4}{3} \times \frac{d(x_5)-3}{2} - \frac{1}{2} - \frac{1}{3} \times \frac{d(x_5)+1}{2} \geq \frac{1}{6}\), a contradiction. Now it remains to consider the case when \(d(x_4) = d(x_6) = 4\). By Lemma 3.5, there is a \(5^+\)-vertex inside \(\{x_2, x_3\}\) that is not dangerous, say \(x_2\), and by Claim 4.2, \(f\) receives \(\frac{1}{6}\) from \(x_2, x_3\) altogether, a contradiction. \(\square\)

Now we need to verify that \(ch_2(f) \geq 0\). If there are at least two \(7^+\)-faces among \(\{f_1, f_2, f_3\}\), then \(ch_2(f) \geq 0\) by R7. Hence we assume at least two of them are 6-face. By
the assumption that \( x_1, w_1, v_1 \) are vice 4-vertices. If \( y_1 \) is not vice and \( f_2 \) is \( 7^+ \)-face, then we are done. Otherwise, if \( y_1 \) is vice, then by Claim 4.4 \( f \) receives at least \( \frac{1}{6} \) from a vertex in \( \{ v_2, v_3, w_2, w_3 \} \). Meanwhile, if \( f_2 \) is a 6-face, then by Claim 4.5 \( f \) receives at least \( \frac{1}{6} \) from a vertex in \( \{ x_2, x_3, x_4, x_5 \} \). This completes the proof.

4.2 Final analysis

Now we shall verify that \( ch_2(x) \geq 0 \) for all \( x \in V(G') \cup F(G') \). It is easy to observe that for each \( v \in V(G') \), if \( f_3(v) < \lfloor \frac{d(v)}{2} \rfloor \), then \( ch_2(v) \geq 0 \). From now on, it suffices to only care the case \( f_3(v) = \lfloor \frac{d(v)}{2} \rfloor \). We assume that each bad vertex satisfies \( d(v) \geq 5 \). For \( v \in V(G') \), denote by \( N_b(v) \) the set of bad neighbors of \( v \) and let \( |N_b(v)| = n_b(v) \).

**Claim 4.6.** Let \( v \) be a vertex of \( G' \) with \( 5 \leq d(v) \leq 12 \). If \( v \) is not dangerous, then \( ch_2(v) \geq 0 \).

**Proof.** If \( v \) is not dangerous, it follows that \( f_{3,3}(v) \leq \lfloor \frac{d(v)-3}{2} \rfloor - 1 \). Accordingly,

\[
ch_2(v) \geq d(v) - 2 - \frac{4}{3}f_{3,3}(v) - f_{3,4}(v) - f_{3b}(v) - \frac{2}{3}f_{4,4}(v) - \frac{5}{6}f_{4b}(v) - f_{6b}(v) \\
- \frac{1}{6}f_{3,4}(v) - \frac{1}{6}f_{3b}(v) - 2 \times \frac{1}{6}f_{4,4}(v) - \frac{1}{6}f_{4b}(v) - \frac{1}{3}(d(v) - f_3(v)) \\
\geq \frac{2}{3}d(v) - \frac{2}{3}f_3(v) - \frac{1}{3}f_{3,3}(v) - \frac{1}{6}(f_3(v) - f_{3,3}(v) - 1) - 2 \\
\geq \frac{d(v) - 8}{6}.
\]

(*)

Thus, \( ch_2(v) \geq 0 \) when \( d(v) \geq 8 \).

In particular, when \( d(v) = 7 \), we get \( ch_2(v) \geq \frac{2}{3} \times 7 - \frac{2}{3} \times 3 - \frac{1}{3} \times 1 - \frac{1}{6} \times 1 - 2 = \frac{1}{6} > 0 \).

When \( d(v) = 6 \). Let \( f_1 = v_1v_2v \), \( f_2 = v_3v_4v \) and \( f_3 = v_5v_6v \) be three 3-faces incident with \( v \), respectively. We discuss the following two cases depending on whether \( v \) is bad:

- **v** is bad but not dangerous in \( G' \);

We may assume that \( f_2 \) and \( f_3 \) are worse, and it follows from \( \{(a), (d), (d)\} \) that \( d(v_1) \geq 4 \), \( d(v_2) \geq 4 \). If \( d(v_1) = d(v_2) = 4 \), then \( \varpi(v) \leq 2 \) by \( \{(e), (e), (k)\} \) and Lemma 3.8, thus \( ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - 2 \times \frac{1}{3} - 1 - \frac{1}{6} = 0 \) by R3 and R8. If \( d(v_1) = 4 \), \( d(v_2) \geq 5 \), then by \( \{(d), (d), (i)\} \), we get \( n_b(v) = 0 \) and \( \varpi(v) \leq 3 \) by Lemma 3.8. Thus \( ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - 2 \times \frac{1}{3} - 3 \times \frac{1}{6} = 0 \) by R5.1. If \( d(v_1) \geq 5 \), \( d(v_2) \geq 5 \), then \( n_b(v) \leq 1 \) by \( \{(d), (d), (h)\} \). If there exists a bad vertex, then it can not be dangerous by Lemma 3.7(2), thus \( ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - 2 \times \frac{1}{3} - 3 \times \frac{1}{6} = 0 \) by R6.3.

- **v** is neither bad nor dangerous in \( G' \);
Case 1. $v$ is incident with a worse face;

W.l.o.g, let $f_3$ be worse and $N_1(v) = N(v) \setminus \{v_5, v_6\}$. Suppose there are two 3-vertices in $N_1(v)$, say $v_1$ and $v_3$, then $d(v_2) \geq 5$, $d(v_4) \geq 5$. By $\{(a), (d), (f)\}$, we obtain $n_b(v) = 0$, by Lemma $3.6$ we get that both $v_2$ and $v_4$ are well vertices, thus $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - 2 \times 2/3 - 2 \times 1/6 = 0$ by R8. Suppose there are only one 3-vertex in $N_1(v)$, say $d(v_2) = 3$, then $d(v_1) \geq 5$. If $d(v_3) = d(v_4) = 4$, we first consider $v_1$ is bad, then $\varpi(v) \leq 2$ by $\{(e), (g), (k)\}, \{(e), (f), (l)\}, \{(d), (g), (l)\}$, thus $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - 2 \times 2/3 - 2 \times 1/6 = 0$ by R3 and R4. Otherwise, $v_1$ is not bad and it follows from $\{(b), (d), (l)\}$ that $\varpi(v) \leq 3$, thus $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - 2 \times 2/3 - 3 \times 1/6 = 0$ by R4. If there exists a 5+-vertex in $\{v_3, v_4\}$, say $v_3$, it follows from $\{(d), (f), (i)\}$ that $n_b(v) \leq 1$. If $v_1$ is bad, then $\varpi(v) \leq 2$ by $\{(e), (g), (m)\}$. On the other hand, if $v_3$ is bad, we can also get $\varpi(v) \leq 2$ by $\{(b), (d), (j)\}$. If there is no bad vertex in $N_1(v)$, then $\varpi(v) \leq 3$ by Lemma $3.8$. Thus $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - \max\{1 + 2 \times 1/3, 2 \times 2 \times 1/3, 5/6 + 2 \times 1/6\} = 0$ by R4 and R5. If $d(v_3) \geq 5$, $d(v_4) \geq 5$, then it follows from $\{(a), (d), (h)\}$ that $v_3$ and $v_4$ are bad at the same time, thus $n_b(v) \leq 2$. If $n_b(v) = 2$, then by Lemma $3.7(2)$, some $v_i$ for $i \in \{3, 4\}$ is bad but not dangerous, and it follows that $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - 2 \times 2/3 - 2 \times 1/6 = 0$ by R4 and R6. If $n_b(v) \leq 1$, then $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - \max\{1 + 2 \times 1/3, 2 \times 2 \times 1/3, 5/6 + 2 \times 1/6\} = 0$ by R4, R6, and R8. Next, we consider $d(z) \geq 4$ for all $z \in N_1(v)$. If $d(z) = 4$ for all $z \not\in N_1(v)$, then $\varpi(v) \leq 4$ by $\{(d), (l), (l)\}$, thus $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - 2 \times 2/3 - 4 \times 1/6 = 0$ by R3. Suppose there exists a 5+-vertex in $N_1(v)$, say $v_1$. If $v_1$ is bad, then $\varpi(v) \leq 3$ by $\{(d), (j), (l)\}$, thus $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - 2 \times 2/3 - 4 \times 1/6 = 0$ by R5.1. If $v_1$ is not bad, then $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - 2 \times 2/3 - 4 \times 1/6 = 0$ by R5.2. Suppose there exists two 5+-vertices in $N_1(v)$, say $v_1$ and $v_3$ or $v_1$ and $v_4$. If $n_b(v) = 2$, then $\varpi(v) \leq 2$ by $\{(d), (j), (j)\}$ or $\{(d), (h), (l)\}$, thus $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - \max\{2 \times 5/6, 1 + 1/2\} = 2 \times 1/6 = 0$ by R3, R5 and R6. If $n_b(v) \leq 1$, then $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - \max\{1/3 + 2/3, 2/3 + 5/6\} - 3 \times 1/6 = 0$ by R5 and R6. If there exists three 5+-vertices in $N_1(v)$, then it follows from $\{(d), (h), (i)\}$ that $n_b(v) \leq 2$ and $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - \max\{1 + 2 \times 2 \times 1/3, 5/6 + 3/4\} - 2 \times 1/6 = 0$ by R5 and R6. If $d(z) \geq 5$ for all $z \not\in N_1(v)$, then it follows from $\{(d), (h), (h)\}$ that $n_b(v) \leq 3$ and $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 1 - 3/4 - 1 - 1/6 = 1/12 > 0$ by R6.

Case 2. $v$ is not incident with a worse face;

If $n_3(v) = 3$, then $n_b(v) \leq 1$ by $\{(a), (f), (f)\}$. In particular, $\varpi(v) \leq 2$ when $n_b(v) = 1$ by Lemma $3.2(6)$, thus $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - \max\{1 + 2 \times 5/6, 2 \times 1/6, 3 \times 5/6 + 3 \times 1/6\} = 0$ by R4. If $n_3(v) = 2$, w.l.o.g, we say $d(v_1) = d(v_3) = 3$, then $d(v_2) \geq 5$, $d(v_4) \geq 5$. Suppose $d(v_5) = d(v_6) = 4$, if $n_b(v) = 2$, then $\varpi(v) \leq 2$ by $\{(g), (g), (k)\}$, and it follows that $ch_2(v) \geq 6 - 2 - 3 \times 1/3 - 2 \times 1 - 3/4 - 2 \times 1/6 = 0$ by R3 and R4. If $n_b(v) \leq 1$, then
Let $d(v_i) = v_1v_2v, f_2 = v_3v_4v$ be two 3-faces incident with $v$. Similarly, we consider whether $v$ is bad.

- $v$ is bad but not dangerous in $G'$.

Assume $f_2$ is worse, if $d(v_i) = 3$, then $d(v_i) \geq 5$, and note that $v_2$ is not bad by $\{(a), (b), (d)\}$, it follows from $\{(a), (b), (d)\}$ that $\varpi(v_i) \leq 2$, thus $\varpi(v_i) = 5 - 2 - 3 \times \frac{1}{3} - 1 - \frac{2}{3} - 2 \times \frac{1}{6} = 0$. If $n_5(v_i) = 3$, then $n_5(v_i) = 0$, i.e. $d(z) = 4$ for all $z \in N(v_i)$, then $d(v_i) \geq 6 - 2 - 3 \times \frac{1}{3} - 3 \times \frac{2}{3} - 6 \times \frac{1}{6} = 0$ by R3. If $n_5(v_i) = 4$, then $d(v_i) \geq 6 - 2 - 3 \times \frac{1}{3} - 3 \times \frac{2}{3} - 6 \times \frac{1}{6} = 0$ by R3. If $n_5(v_i) = 5$, then $d(v_i) \geq 6 - 2 - 3 \times \frac{1}{3} - 3 \times \frac{2}{3} - 6 \times \frac{1}{6} = 0$ by R3. If $n_5(v_i) = 6$, then $d(v_i) \geq 6 - 2 - 3 \times \frac{1}{3} - 3 \times \frac{2}{3} - 6 \times \frac{1}{6} = 0$ by R3. If $n_5(v_i) = 7$, then $d(v_i) \geq 6 - 2 - 3 \times \frac{1}{3} - 3 \times \frac{2}{3} - 6 \times \frac{1}{6} = 0$ by R3. If $n_5(v_i) = 8$, then $d(v_i) \geq 6 - 2 - 3 \times \frac{1}{3} - 3 \times \frac{2}{3} - 6 \times \frac{1}{6} = 0$ by R3.
by R3 and R4. If \( d(v_1) = d(v_2) = 4 \), then \( \omega(v) \leq 2 \) by \{ (a), (d), (k) \}, thus \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 1 - \frac{2}{3} - 2 \times \frac{1}{6} = 0 \) by R3. If \( d(v_1) = 4 \) and \( d(v_2) \geq 5 \), then \( v_2 \) is not bad by \{ (d), (i) \} and \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 1 - \frac{2}{3} - 3 \times \frac{1}{6} = 0 \) by R3 and R5. If \( d(v_1) \geq 5 \) and \( d(v_2) \geq 5 \), then it follows that \( n_b(v) \leq 1 \) by \{ (d), (h) \}. If \( n_b(v) = 1 \), w.l.o.g., let \( v_1 \) be a bad vertex. If \( v_1 \) is also dangerous, then \( d(v_5) \geq 4 \) by Lemma 3.7(1), and it follows that \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 1 - \frac{2}{3} - 2 \times \frac{1}{6} = 0 \). If \( n_b(v) = 0 \), then \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 1 - \frac{2}{3} - 2 \times \frac{1}{6} = 0 \).

- \( v \) is neither bad nor dangerous in \( G' \);

Let \( N_3(v) = \{ v_1, v_2, v_3, v_4 \} \), we denote \( N_{3b}(v), N'_3(v) \) the set of bad vertices and 3-vertices in \( N_3(v) \) respectively. For simplicity, let \( n_{3b}(v) = |N_{3b}(v)|, n'_3(v) = |N'_3| \). If \( n'_3(v) = 2 \), say \( v_1 \) and \( v_3 \), then \( d(v_2) \geq 5 \). If \( d(v_3) = 3 \), then \( d(v_4) \geq 5 \). We get \( n_{3b}(v) \leq 1 \) by \{ (f), (f) \}. If \( n_{3b}(v) = 1 \), then \( \omega(v) \leq 1 \) by Lemma 3.2(5), thus \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 1 - \frac{2}{3} - 2 \times \frac{1}{6} = 0 \). Otherwise, \( n_{3b}(v) = 0 \), if \( d(v_5) \geq 4 \), then \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 2 \times \frac{1}{6} = 0 \). If \( d(v_5) = 3 \), it follows from Lemma 3.6 that both \( v_2 \) and \( v_4 \) are well vertices, thus \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 2 \times \frac{1}{6} - 2 \times \frac{1}{6} = 0 \). If \( n'_3(v) = 1 \), say \( v_1 \), then \( d(v_2) \geq 5 \). When \( d(v_3) = d(v_4) = 4 \), if \( v_1 \) is bad, then \( \omega(v) \leq 2 \) by \{ (a), (g), (k) \}, thus \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 1 - \frac{2}{3} - 2 \times \frac{1}{6} = 0 \) by R3 and R4. Otherwise, \( v_1 \) is not bad, then \( \omega(v) \leq 3 \) by \{ (a), (b), (l) \}, thus \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - \frac{5}{6} - \frac{2}{3} - 3 \times \frac{1}{6} = 0 \). When \( d(v_i) \geq 5 \) for some \( i \in \{ 3, 4 \} \), if \( n_{3b}(v) = 2 \), then \( d(v_5) \geq 4 \) by \{ (a), (f), (i) \}. Moreover, \( \omega(v) \leq 1 \) by \{ (g), (j) \}, thus \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 1 - \frac{2}{3} - 2 \times \frac{1}{6} = 0 \). If \( n_{3b}(v) = 1 \), first, suppose \( v_i \) is bad, then \( \omega(v) \leq 2 \) by \{ (a), (b), (i) \}, thus \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 2 \times \frac{1}{6} = 0 \). Second, suppose \( v_2 \) is bad, it follows from \{ (a), (g), (m) \} that \( \omega(v) \leq 2 \), thus \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 1 - \frac{2}{3} - 2 \times \frac{1}{6} = 0 \). If neither \( v_2 \) nor \( v_i \) is bad, then \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 1 - \frac{2}{3} - 2 \times \frac{1}{6} = 0 \). When \( d(v_3) \geq 5 \), \( d(v_4) \geq 5 \), \( n_{3b}(v) \leq 2 \) by \{ (f), (h) \}. If \( n_{3b}(v) = 2 \), we get \( \omega(v) \leq 1 \) by Lemma 3.7(1), \{ (a), (b), (h) \}, thus \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - \frac{5}{6} - \frac{2}{3} - 3 \times \frac{1}{6} = 0 \). If \( d(v_5) = 3 \) and \( n_{3b}(v) = 2 \), we get the \( \omega(v) \leq 2 \) by \{ (a), (h), (l) \}. If \( d(v_5) = 3 \), then \( \omega(v) \leq 2 \) by \{ (a), (m), (m) \}, it follows that \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 2 \times \frac{1}{6} - 4 \times \frac{1}{6} = 0 \) by R3 and R5. If there exist a 5\textsuperscript{+} vertex in \( N_3(v) \), say \( v_1 \), if \( v_1 \) is bad, then \( \omega(v) \leq \frac{1}{3} \) by \{ (a), (j), (l) \}, thus \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 2 \times \frac{1}{6} - 4 \times \frac{1}{6} = 0 \). Otherwise, \( v_1 \) is not bad, it follows that \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 2 \times \frac{1}{6} - 4 \times \frac{1}{6} = 0 \) by R3 and R5. If there are two 5\textsuperscript{+}-vertices in \( N_3(v) \), say \( v_1 \) and \( v_2 \) or \( v_1 \) and \( v_3 \), if \( n_{3b}(v) = 2 \), then \( \omega(v) \leq 2 \) by \{ (a), (h), (l) \}, \{ (a), (j), (j) \}, then \( ch_2(v) = 5 - 2 - 3 \times \frac{1}{3} - \max \{ 2 \times \frac{5}{6}, 1 + \frac{2}{3} \} - 2 \times \frac{1}{6} = 0 \). Otherwise if \( n_{3b}(v) \leq 1 \), then \( ch_2(v) = 5 - 2 - 3 \times \frac{1}{3} - \max \{ \frac{5}{6} + \frac{2}{3}, \frac{3}{4} + \frac{2}{3} \} - 3 \times \frac{1}{6} = 0 \). If there are three 5\textsuperscript{+}-vertices in \( N_3(v) \), and if \( n_{3b}(v) = 3 \), then \( \omega(v) \leq 1 \) by \{ (a), (h), (j) \}, thus
\( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - \frac{5}{6} - 1 - \frac{1}{6} = 0 \) by R5 and R6. Otherwise \( n_{3b}(v) \leq 2 \), then \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - \max\{\frac{2}{3}, 1\} = 2 \times \frac{1}{6} = 0 \). If \( d(z) \geq 5 \) for all \( z \in N_2(v) \), then \( n_{3b}(v) \leq 3 \) by \{\( (h), (h) \}\), thus \( ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - 1 - \frac{3}{4} - \frac{1}{6} = \frac{1}{12} \geq 0 \).

This completes the proof of Claim \[4.6\].

Now, we are ready to verify all vertices in \( G' \) satisfying \( ch_2(v) \geq 0 \).

**Let \( v \) be a 3-vertex in \( G' \).** Then \( ch_2(v) = ch_1(v) = 1 - 3 \times \frac{1}{3} = 0 \) by R1.

**Let \( v \) be a 4-vertex in \( G' \).** Then \( ch_2(v) = ch_1(v) \geq 2 - \max\{1 + 3 \times \frac{1}{3}, \frac{1}{2} + \frac{2}{3} + 2 \times \frac{1}{3}, \frac{5}{6} + 3 \times \frac{1}{3}\} = 0 \) by R2.

**Let \( v \) be a 5\(^+\)-vertex in \( G' \).** Suppose \( v \) is not dangerous, then \( ch_2(v) \geq 0 \) by Claim \[4.6\]. Next, we consider the case \( v \) is dangerous.

If \( d(v) \) is even, then \( f_{3,3}(v) = \frac{d(v)}{2} - 2 \), thus we obtain that \( ch_2(v) = d(v) - \frac{4}{3} (\frac{d(v)}{2} - 2) - \frac{1}{3} (d(v) - d_3(v)) - 1 - \frac{7}{6} = \frac{d(v) - 9}{6} \geq 0 \) when \( d(v) \geq 9 \).

**Let \( v \) be a 8-vertex in \( G' \).** Note that \( ch_2(v) = 2 - \frac{7}{6} f_{3,4}(v) - f_{4,4}(v) - f_{4b}(v) \). Let \( f_3 \) as well as \( f_4 \) be worst and \( f_1 \) and \( f_2 \) be the rest two 3-faces. If \( v \) is bad, then \( \varpi(v) \leq 2 \) by \{\( (c), (c), (e), (l) \}\), and it follows that \( ch_2(v) \geq 8 - 2 - 4 \times \frac{1}{3} - 2 \times \frac{4}{3} - 1 - \frac{2}{3} - 2 \times \frac{1}{6} = 0 \) by R3 and R5. Otherwise, \( ch_2(v) \geq 8 - 2 - 4 \times \frac{1}{3} - 2 \times \frac{4}{3} - \max\{2 \times \frac{2}{3} + 4 \times \frac{1}{6}, 2 \times \frac{5}{6} + 2 \times \frac{1}{6}\} = 0 \) by R3-R5.

**Let \( v \) be a 6-vertex in \( G' \).** We assume that \( f_1 = v_1v_2v \), \( f_2 = v_3v_4v \) and \( f_3 \) is worst.

- \( v \) is bad and dangerous in \( G' \);

W.l.o.g, assume \( f_3 \) is worse. If \( d(v_1) = d(v_2) = 4 \), then \( ch_2(v) = ch_1(v) = 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - 1 - \frac{2}{3} = 0 \). If \( d(v_i) \geq 5 \) for some \( i \in \{1, 2\} \), then \( v_i \) can not be bad by \{\( (c), (d), (i) \}\), and \( \varpi(v) \leq 1 \) by \{\( (c), (e), (m) \}\), thus \( ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - 1 - \frac{1}{2} - \frac{1}{6} = 0 \) by R4 and R5. If \( d(v_i) \geq 5 \) for all \( i \in \{1, 2\} \), then \( n_b(v) \leq 1 \) by \{\( (c), (d), (h) \}\), thus \( ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - 1 - \frac{1}{2} - \frac{1}{6} = 0 \) by R3 and R6.

- \( v \) is not bad but dangerous in \( G' \);

**Case 1.** \( n_3(v) = 4 \);

Then \( n_{5^+}(v) = 2 \). It follows that \( ch_2(v) = ch_1(v) = 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - 2 \times \frac{5}{6} = 0 \).

**Case 2.** \( n_3(v) = 3 \);

Let \( v_1 \) be another 3-vertex, then \( d(v_3) \geq 5 \). If \( d(v_3) = d(v_4) = 4 \), then \( n_b(v) = 0 \) by Lemma \[3.2\] \( (2) \), and \( \varpi(v) \leq 1 \) by \{\( (a) \) or \( (b) \), \( (c) \), \( (k) \)\}, thus \( ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{5}{6} - \frac{2}{3} - \frac{1}{6} = 0 \) by R3 and R4. If \( d(v_i) \geq 5 \) for some \( i \in \{3, 4\} \), then \( n_b(v) = 0 \) by \{\( (a), (c), (i) \}\} and Lemma \[3.2\]. Moreover, both \( v_3 \) and \( v_i \) are well vertices by Lemma \[3.6\], thus \( ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{5}{6} - \frac{2}{3} - 2 \times \frac{1}{12} = 0 \) by R8. If \( d(v_i) \geq 5 \) for all \( i \in \{3, 4\} \), note that \( n_b(v) = 1 \) by \{\( (a), (c), (h) \)\}.
and \(v_2\) must be well-vertices by Lemma \[3.6\], then \(ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{5}{6} - \frac{3}{4} - \frac{1}{12} = 0\) by R8.

**Case 3.** \(n_3(v) = 2\);

If \(n_4(v) = 4\), then \(\varpi(v) \leq 2\) by \{(c), (k), (l)\}, it follows that \(ch_2(v) = 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} - 2 \times \frac{1}{6} = 0\) by R3. We denote \(N_4(v) = N(v) \setminus \{v_5, v_6\}\), suppose \(n_5^+(v) = 1\). If \(n_4(v) = 1\), then \(\varpi(v) \leq 1\) by \{(c), (i), (l)\}, \{(c), (j), (k)\}, it follows that \(ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} - 2 \times \frac{5}{6} - \frac{1}{6} = 0\) by R3 and R5. Otherwise if \(n_b(v) = 0\), then \(\varpi(v) \leq 2\) by \{(c), (l), (m)\}, it follows that \(ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} - 2 \times \frac{5}{6} = 0\) by R3 and R5. If \(n_5^+(v) = 2\), say \(v_1, v_2\) or \(v_1, v_3\). In the former case, by Lemma \[3.2\](4), we get that \(n_b(v) \leq 1\). Suppose \(n_6(v) = 1\), say \(v_1\), if \(v_1\) is also dangerous, then \(\varpi(v) \leq 1\) by Lemma \[3.7\](6), then \(ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{3}{4} - \frac{1}{6} = \frac{11}{12} > 0\). Otherwise, we get \(ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} - 2 \times \frac{1}{6} = 0\) by R6. Suppose \(n_b(v) = 0\), then \(ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} - 2 \times \frac{7}{12} = 0\) by R3 and R6. In the latter case, suppose \(n_b(v) = 2\), then \(ch_2(v) = ch_1(v) = 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} - 2 \times \frac{5}{6} = 0\). If \(n_b(v) = 1\), then \(\varpi(v) \leq 1\) by \{(c), (j), (m)\}, it follows that \(ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} - \frac{5}{6} - \frac{1}{6} = 0\) by R5. Otherwise if \(n_b(v) = 0\), it follows that \(ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} - 2 \times \frac{7}{12} = 0\) by R5. If \(n_5^+(v) = 3\), then \(n_b(v) \leq 2\) by \{(c), (h), (i)\}. If \(n_b(v) = 2\), by the same argument, we have \(ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{4}{3} - \max\{\frac{5}{6}, \frac{3}{4}, \frac{2}{3}, \frac{5}{6}, \frac{1}{4}, \frac{1}{4}\} = 0\) by R5 and R6. If \(n_6(v) = 0\), then \(ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} - \frac{4}{3} - \frac{4}{3} - \frac{1}{6} = 0\) by R5 and R6. If \(n_5^+(v) = 4\), then there are at most two bad vertices by Lemma \[3.2\](4), thus \(ch_2(v) \geq 6 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} - \frac{3}{2} = 0\) by R3 and R6.

If \(d(v)\) is odd, note that \(v\) is also bad, it follows that \(f_{3,4}(v) = f_{3b}(v) = f_{4b}(v) = f_{bb}(v) = 0\). Then \(ch_2(v) \geq d(v) - 2 - \frac{4}{3}f_{3,3}(v) - \frac{1}{3}d(v) - f_3(v) = \frac{d(v) - 7}{6} \geq 0\) when \(d(v) \geq 7\).

**Let \(v\) be a 5-vertex in \(G'\).** Let \(f_1 = v_1v_2v\) and assume that \(f_2\) is worst. If there exists a 3-vertex lying on \(f_1\), say \(v_1\), then \(d(v_2) \geq 5\) and \(v_2\) is not bad by \{(c), (f)\}. It follows that \(ch_2(v) = ch_1(v) \geq 5 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} = 0\) by R3 and R4. If \(d(v_1) = d(v_2) = 4\), then \(ch_2(v) = ch_1(v) = 5 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{2}{3} = 0\) by R3. Otherwise, there exists a \(5^+\)-vertex lying on \(f_1\) which is not bad by \{(c), (i)\}, then \(\varpi(v) \leq 1\) by \{(a) or (b), (c), (m)\}. It follows that \(ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{1}{2} - \frac{1}{6} = 0\) by R3 and R5. If there are two \(5^+\)-vertices lying on \(f_1\), it follows from \{(c), (h)\} that \(n_b(v) \leq 1\), thus \(ch_2(v) \geq 5 - 2 - 3 \times \frac{1}{3} - \frac{4}{3} - \frac{1}{2} - \frac{1}{6} = 0\) by R3 and R5.

**Let \(f\) be a 6+-face in \(G'\).** Then \(ch_2(f) \geq 0\) by R1 and R7.

**Let \(f\) be a 3-face in \(G'\).** Let \(f = v_1v_2v_3\), we next consider different cases corresponding to the shape of \(f\). If \(f\) is poor, it follows that \(ch_2(f) \geq -2 + \frac{1}{3} + 2 \times \frac{2}{3} + \min\{\frac{1}{3}, 2 \times \frac{1}{6}\} = 0\) by R1, R7 and R8. In particular, if there exists at least one 4-vertex which is not worse than \(f\), then \(ch_2(f) \geq -2 + 1 + \frac{1}{3} + \frac{2}{3} = 0\) by R2.1. If \(d(v_1) = 3, 3 \leq d(v_2) \leq 4, d(v_3) \geq 5\), note that \(v_2\)
cannot be bad, then $ch_2(f) \geq -2 + \frac{1}{3} + \min \{ \frac{1}{3} + \frac{4}{3}, \frac{2}{3} + 1 \} = 0$ by R3. If $d(v_1) = 3$, $d(v_2) \geq 5$, $d(v_3) \geq 5$, by Lemma 3.2, there is at most one bad vertex contained in $\{v_2, v_3\}$. It follows that $ch_2(f) \geq -2 + \frac{1}{3} + \min \{ \frac{2}{3} + \frac{5}{6}, 2 \times \frac{2}{3} \} = 0$ by R4. If $d(v_i) = 4$ for each $i \in \{1, 2, 3\}$, then $ch_2(f) = -2 + 3 \times \frac{2}{3} = 0$ by R2.3. If $d(v_1) = 4, d(v_2) = 4, d(v_3) \geq 5$, it follows that there is at most one bad vertex contained in $\{v_1, v_2, v_3\}$, then $ch_2(f) \geq -2 + \frac{1}{3} + \min \{ \frac{1}{3} + \frac{4}{3}, \frac{2}{3} + 1 \} = 0$ by R6.

Hence, $ch_2(x) \geq 0$ for all $x \in V(G') \cup F(G')$, this contradiction completes the proof of Theorem 1.

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