Exact-2-Relation Graphs

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Abstract

Pairwise compatibility graphs (PCGs) with non-negative integer edge weights recently have been used to describe rare evolutionary events and scenarios with horizontal gene transfer. Here we consider the case that vertices are separated by exactly two discrete events: Given a tree $T$ with leaf set $L$ and edge-weights $\lambda : E(T) \rightarrow \mathbb{N}_0$, the non-negative integer pairwise compatibility graph $\text{nniPCG}(T, \lambda, 2, 2)$ has vertex set $L$ and $xy$ is an edge whenever the sum of the non-negative integer weights along the unique path from $x$ to $y$ in $T$ equals 2. A graph $G$ has a representation as $\text{nniPCG}(T, \lambda, 2, 2)$ if and only if its point-determining quotient $G/\sim$ is a block graph, where two vertices are in relation $\sim$ if they have the same neighborhood in $G$. If $G$ is of this type, a labeled tree $(T, \lambda)$ explaining $G$ can be constructed efficiently. In addition, we consider an oriented version of this class of graphs.

Keywords: Pairwise compatibility graphs, edge-labeled trees, thin graphs, block graphs, oriented graphs

2000 MSC: 05C75, 05C05, 92B10

1. Introduction

Consider a tree $T$ with leaf set $V$ and a non-negative edge-weight function $\ell : E(T) \rightarrow \mathbb{R}_0^+$. Denote by $P(x, y)$ the unique path between $x$ and $y$ in $T$. The canonical distance function $d_{T,\ell} : V \times V \rightarrow \mathbb{R}_0^+$ is then defined by

$$d_{T,\ell}(x, y) = \sum_{e \in P(x, y)} \ell(e)$$

(1)

This definition is the starting point for mathematical phylogenetics, which is centered around finite additive metric spaces and their generalizations [1, 2]. It also serves as a basis for defining a large class of graphs that in the recent past has received considerable attention. The pairwise compatibility graphs $G = \text{PCG}(T, \ell, d_{\text{min}}, d_{\text{max}})$ has vertex set $V$ and edges

$$uv \in E(G) \quad \text{if and only if} \quad d_{\text{min}} \leq d_{T,\ell}(u, v) \leq d_{\text{max}}$$

(2)

Originally introduced in the context of phylogenetics [3], they have received considerable interest in the last years, see [4, 5] and the references therein, as well as [6, 7, 8]. A further generalization of “multi-interval” PCGs in explored in [9]. In the setting of PCGs and most phylogenetic applications one usually stipulates that $\ell(e) > 0$, measuring e.g. the time between two distinct events associated with adjacent vertices of $T$. A class of graphs that is conceptually closely are the exact $k$-leaf powers [10], for which $A(e) = 1$ for all edges of $T$ and $d_{\text{min}} = d_{\text{max}} = k$.

In an alternative interpretation, $\ell(e)$ models the number of discrete evolutionary events along an edge of $e$ of $T$. This is of interest in particular in the context of so-called rare genomic changes such as the gene or loss of a particular

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gene of gene family or a particular genomic rearrangement [11]. Some of these convey phylogenetic information that is (nearly) free of homoplasy, i.e., the independent occurrence in independent lineages. Examples of such rare events are the emergence of novel microRNA families [12] or rearrangements of the genomic gene order [13]. Since such events are very unlikely to have occurred more than once in the same manner, they identify phylogenetic groupings that share such an innovation with very little ambiguity. This provides information to resolve also parts of the phylogenetic tree where classical, sequence-based methods fail [14]. In this context it is necessary to allow \( \ell(e) = 0 \) because the events of interest are by definition so rare that not all taxa will be distinguished by them. In the same vein it important that events are discrete and hence an integer-valued weight function \( \lambda : E(T) \to \mathbb{N}_0 \). Both conditions on \( \ell \) cause subtle but important differences in comparison with the usual definition of PCGs that requires non-zero edge length but otherwise allows arbitrary real values. We will denote these “non-negative integer pairwise compatibility graph” by \( \text{nniPCG} \) to distinguish them from the better studied class of PCGs with non-zero real-valued edge weights \( \lambda \).

The special case in which two leaves \( x \) and \( y \) of \( T \) are separated by a single event, corresponding to graphs of the form \( \text{nniPCG}(T,\lambda,1,1) \), was explored in [15] as a model of rare events in evolution. It turns out that this graph class coincides with the forests. The graphs \( \text{nniPCG}(T,\lambda,1,\infty) \) requiring at least one event along the path between two leaves also have a very simple structure: they are exactly the complete multipartite graphs [16].

Considering a rooted tree \( \bar{T} \) instead of an unrooted tree \( T \), it is natural to consider the digraphs with edges \((x,y)\) whenever a certain number of events occurred between the last common ancestor \( \text{lca}(x,y) \) and \( y \). This construction appears naturally in the context of horizontal gene transfer (HGT), where one asks whether \( \text{lca}(x,y) \) and \( y \) are separated by at least one HGT event\(^1\). This gives rise to the class of Fitch graphs [17], which form a subclass of di-cographs introduced by [18]. Their underlying undirected graphs are exactly the \( \text{nniPCG}(T,\lambda,1,\infty) \), i.e., the complete multipartite graphs.

A related construction requires a certain number of events between \( \text{lca}(x,y) \) and \( y \) and excludes all events between \( \text{lca}(x,y) \) and \( x \). This class of graphs appears naturally when events are directed, i.e., when it is (in general) no possible to revert the effect of an operation in a single step. Probably the best studied type of single genomic events of this type are so-called Tandem-Duplication-Random-Loss events, during which a genomic interval is duplicated and then one of the two copies of each gene is lost at random [19]. The antisymmetric digraphs obtained by single events are characterized in [15].

Our interest in the graphs \( \text{nniPCG}(T,\lambda,k,k) \) for general \( k \geq 1 \) also stems from rare-event phylogenetic data. Since we assume an underlying tree, the distance matrix \( d_T \) is additive and its entries are small non-negative integers. The fact that all edge lengths \( \ell(e) \) are also integers of course imposes additional constraints. As demonstrated e.g. in the context of orthology assignment (a related problem with vertex labeled trees for which the corresponding graphs turn out to be cographs), graph editing can be employed to correct empirically estimated input graphs [20]. This approach requires, however, that constraint on the graphs that can appear are known. In the case of rare-event phylogenetics, we know that the graph with edge set \( \{xy|d_T(x,y) = k\} \) must be a \( \text{nniPCG}(T,\lambda,k,k) \). In the rare-event scenario, the number of pairs of nodes with \( d_T(x,y) = k \) will quickly decrease with \( k \), so that the empirical input graphs will have few edges for larger values of \( k \) and thus rarely reveal obstructions. Hence only small value of \( k \) are of practical value for detecting measurement errors in the data. Since the \( \text{nniPCG}(T,\lambda,1,1) \) are forests, the corresponding graph editing problem amounts to identifying spanning forests, and possible false positive events are edges in cycles. False negatives are not detectable for \( k = 1 \) since there are no non-tree graphs that would become trees by inserting edges. They could be detected, however, as missing edges in the empirical graph for \( \text{lca}(x,y) \) and \( y \) and \( k \) compared to the most similar member of \( \text{nniPCG}(T,\lambda,2,2) \). In this contribution, therefore, we are interested in the characterization of the graphs \( \text{nniPCG}(T,\lambda,2,2) \), in which edges correspond to exactly two events between two leaves. This graph class is very different from the exact-2-leaf power graphs, which are known to coincide with the disjoint unions of cliques [21, 10]. In contrast, we shall see below that e.g. every path also has a representation as \( \text{nniPCG}(T,\lambda,2,2) \).

This contribution is organized as follows: We first consider a few general properties of the slightly more general \( \text{exactly-k-relation} \) before investigating for some small graphs and simple graph families whether they can be represented with respect to the \( \text{exactly-2-relation} \) on the leaf set of some tree. Here, we consider the case that \( d_T(x,y) > 0 \), i.e., that all leaves are separated by at least one event, and then relax this constraint and characterize the entire graph class \( \text{nniPCG}(T,\lambda,2,2) \) for non-negative, integer \( \lambda \). Our main result is that these graphs are those whose

\(^1\)HGT refer to the import of gene from an unrelated species e.g. through an infection, ingestion, acquisition via a plasmid.
Furthermore, we say (no information on position of root. In fact, it is well known [23, 24] that a metric $d$ is Euclidean if and only if $d_{min} = d_{max}$. This is also known as the exact-$k$-relation, which is different from the exact-$k$-leaf power graphs studies systematically in [10]. While it is true that every PCG($T, \lambda, d_{min}, d_{max}$) with non-negative real weights $\lambda$ and bounds $d_{min}$ and $d_{max}$ also has a representation as PCG($T, \lambda, d_{min}, d_{max}$) with integer weights and bounds [22, Lemma 2], the restriction to integer weights clearly a restriction to integer edge weights on the one hand, and the admission of zero-weights on the other hand, make the graphs non-niPCG($T, \lambda, k$) quite different from the exact-$k$-leaf power graphs studies systematically in [10].

Definition 1. Let $(T, \lambda)$ be an unrooted tree with leaf set $L$ and edge-labeling function $\lambda : E(T) \rightarrow \mathbb{N}_0$. For $x, y \in L$ we consider the exactly-$k$-relation $\sim^k$ defined by $x \sim^k y$ if the (unique) path $(x, y)$ from $x$ to $y$ in $T$ satisfies $\sum_{e \in P(x,y)} \lambda(e) = k$. Furthermore, we say $(T, \lambda)$ explains a graph $G(L, E)$ (with respect to the exactly-$k$-relation) if $(x, y) \in E$ if and only if $x \sim^k y$.

We consider unrooted instead of rooted tree since the distances $d_T(x, y)$ and thus the exactly-$k$-relation $\sim^k$ contains no information on position of root. In fact, it is well known [23, 24] that a metric $d$ of the form (1) uniquely defines an unrooted tree. Therefore, one can only hope to reconstruct the unrooted tree $T$.

Definition 2. The edge-labeled tree $(T, \lambda)$ displays the edge labeled tree $(T', \lambda')$ if $(T', \lambda')$ can be obtained from $(T, \lambda)$ by first removing every edge and vertex from $(T, \lambda)$ that is not contained in a path connecting two leaves of $T'$, and then contracting every path $P(u, v)$ in the remainder of $T$ that has only interior vertices of degree 2 by a single edge $e'$ in $T'$ with label $\lambda'(e') = \sum_{e \in P(u,v)} \lambda(e)$.

In particular, therefore it is sufficient to consider phylogenetic trees, that is, trees $T$ in which every interior node $x \in V(T) \setminus L$ has degree at least 3. Fig. 1 gives an example. A simple, but important consequence of Definition 1 is the following

Lemma 3. If $(T, \lambda)$ displays $(T', \lambda')$, $(T, \lambda)$ explains $G(L, \sim^k)$ and $(T', \lambda')$ explains $G'(L', \sim^{k'})$. Then $G'(L', \sim^{k'}) = G(L, \sim^k)[L']$, the subgraph of $G(L, \sim^k)$ induced by $L'$.

Proof. If $(T, \lambda)$ displays $(T', \lambda')$ then $\sum_{e \in P_T(u,v)} \lambda(e) = \sum_{e \in P_{T'}(u,v)} \lambda'(e)$ for all $u, v \in L(T') \subseteq L(T)$, and thus we conclude that for all $u, v \in L(T')$, we have $u \sim_{(T', \lambda')} v$ if and only if $u \sim_{(T, \lambda)} v$, i.e., $G'(L', \sim^{k'})$ is the subgraph of $G(L, \sim^k)$ induced by $L'$. \qed
Figure 2: The false twin classes are indicated by dashed outlines in the graph \( G \) (middle) and form the vertices (in red) of the quotient graph \( G/\sim \) on the right. For comparison, the orginal edges between the \( \sim \) classes are shown. The subgraph \( G[r_1, r_2, r_3, r_4, r_5] \) of \( G \) induced by one representative \( r_i \) for each \( \sim \)-class is shown on the left. It is isomorphic to \( G/\sim \) irrespective of the choice of the representative vertices \( r_i \) for the \( \sim \) classes.

It follows that “being explained with respect to the exactly-\( k \)-relation” is a hereditary graph property for all \( k \).

We also note the following immediate consequence of the definition.

**Lemma 4.** If \((T, \lambda)\) explains \( G \) with respect to \( \lambda \), then \((T, kl)\) explains \( G \) with respect to \( k \).

**Lemma 5.** Let \( G \) be a graph with connected components \( G_i, \ i = 1, \ldots, N \). Then there is an edge-labeled tree \((T, \lambda)\) explaining \( T \) with respect to \( k \) if and only if there are edge labeled trees \((T_i, \lambda_i)\) explaining \( G_i \) for all \( i = 1, \ldots, N \).

**Proof.** The condition is necessary because of heredity. In order to see sufficiency, we can construct \((T, \lambda)\) from the disjoint union of the \((T_i, \lambda_i)\) in the following way: first we arrange them as an arbitrary tree \( T \). Then we replace each \((K_2, \lambda(e) = k)\) by \( S_2 \) with the two edges \( e' \) and \( e'' \) labeled such that \( \lambda(e') + \lambda(e'') = k \). Now choose for each tree \( T_i \) an arbitrary inner vertex \( x_i \) in \( T_i \neq K_1 \) and the unique vertex \( x_i \) if \( T_i = K_1 \). Finally, we connect \( x_i \) and \( x_j \) by an edge \( e_{ij} \) with \( \lambda(e_{ij}) = k + 1 \) if and only if \( T_i \) and \( T_j \) are adjacent in \( T \). To verify that \((T, \lambda)\) indeed explains \( G \) we observe: (i) If \( x \) and \( y \) are leaves from different connected components of \( G \), they are located in different subtrees \( T \) and thus the path connecting them contains one of the edges label \( k + 1 \), thus \( x \) and \( y \) are not in relation \( \sim \).

It is therefore sufficient to consider connected graphs.

**Definition 6.** An edge-labeled graph \((T, \lambda)\) is canonical if \( T \) is phylogenetic and \( \lambda(e) \neq 0 \) for all interior edges.

**Lemma 7.** Let \((\hat{T}, \hat{\lambda})\) be the edge labeled tree obtained from \((T, \lambda)\) by (1) replacing every path \( P(u, v) \) in \( T \) whose interior vertices have degree 2 by a single edge \( e' \) in \( T' \) with label \( \lambda(e') = \sum_{u \in P(u, v)} \lambda(e) \) and (2) contracting every interior edge with \( \lambda(e) = 0 \). The tree \((\hat{T}, \hat{\lambda})\) is uniquely defined, canonical, and explains the same graph as \((T, \lambda)\).

**Proof.** The maximal paths with interior vertices of degree 2 in \( T \) are disjoint and thus can be treated independently. By construction, any such path \( P \) can also be stepwisely replaced by edges, eventually arriving at the same edge weight for the single edge that remains. Given \( T \), the resulting tree \( \hat{T} \) is therefore unique and contains no vertex of degree 2. It is therefore phylogenetic. Since an interior edge with label \( \lambda(e) = 0 \) does not contribute the total weight of any path that runs through it, it can be contracted without changing the total path weights between leaves. Thus \((T, \lambda)\) and \((\hat{T}, \hat{\lambda})\) explain the same graph.

**Definition 8.** Let \( G \) be a graph. For each \( x \in V(G) \) denote by \( N(x) \) the neighbors of \( x \). Two vertices \( x \) and \( y \) are false twins, \( x \not\sim y \), if \( N(x) = N(y) \).

In contrast, true twins, which play no role here, satisfy \( N(x) \cup \{x\} = N(y) \cup \{y\} \). By definition, false twins \( x \not\sim y \) are non-adjacent, while true twins are always adjacent [25].

The false twin (R-thinness) relation \( \not\sim \) has been well studied in the literature, in particular in the context of graph products [26]. It is well known that \( \not\sim \) is an equivalence relation, see e.g. [26, sect. 8.2]. Its equivalence classes,
which we denote by \( R_i, i = 1, \ldots, h \), are totally disconnected in \( G \) because, by definition, \( x \notin N(x) \). Denote by 
\[ G[r_1, r_2, \ldots, r_h] \]
the subgraph of \( G \) induced by one arbitrarily chosen representative \( r_i \in R_i \) of each false twin class. Since for any \( x \in R_i \) and \( y \in R_j \) we have \( xy \in E(G) \) if and only \( x' y' \in E(G) \) for all \( x' \in R_i \) and all \( y' \in R_j \) we observe that 
\[ G[r_1, r_2, \ldots, r_h] \]
and the quotient graph \( G/\sim \) are isomorphic. An illustration is given in Fig. 2.

**Lemma 9.** Let \((T, \lambda)\) be a canonical tree explaining a connected graph \( G \) with respect to \( \sim \), and let \( W \) be a set of sibling leaves attached to the same parent \( q \) with \( \lambda(qw) = 0 \) for all \( w \in W \). Then \( W \) is contained in a false twin class for the graph \( G \) explained by \((T, \lambda)\) with respect to \( \sim \) for all \( k > 0 \).

**Proof.** Consider a node \( y \in L \setminus W \). Then the total weight of the path between \( y \) and every \( w \in W \) is the same. Furthermore, the total path weight between any two vertices in \( w', w'' \in W \) is 0, i.e., there is no edge between \( w' \) and \( w'' \). Thus \( N(w') = N(w'') \), i.e., \( w' \sim w'' \) for all \( w', w'' \in W \).

A graph is called \textit{R-thin} \cite{27}, \textit{point determining graph} \cite{28} or \textit{mating graph} \cite{29} if \( \sim \) is discrete, i.e., every false twin class consists of only a single point. Clearly, \( G/\sim \) is R-thin. R-thin graphs have also been studied from the point of view of combinatorial enumeration \cite{30,31}. Algorithms for prime-factorization of graphs, furthermore, often operate on \( G/\sim \), since R-thinness ensures uniqueness of the factorization and allows for highly efficient algorithms \cite{27,32,26}. Below we show that it also suffices to consider \( G/\sim \), i.e., R-thin graphs, in our setting. Indeed, a simple consequence of Lemma 9 is

**Corollary 10.** If \( G \) is R-thin and \((T, \lambda)\) is a canonical tree explaining \( G \) with respect to \( \sim \), then \( \sim \) is discrete.

**Algorithm 1** Compute \((T^*, \lambda^*)\) from \((T^*, \lambda^*)\) and false twin classes \( R_i \) with representatives \( r_i \in R_i \).

**Require:** \((T^*, \lambda^*), (r_i, R_i)\) for \( i = 1, \ldots, h \).

1. **for all** false twin classes \( R_i \) with \(|R_i| \geq 1\) **do**
2. \quad \( q \leftarrow \) unique neighbor of leaf \( r_i \) in \((T^*, \lambda^*)\)
3. \quad **remove** \( r_i \) from \((T^*, \lambda^*)\)
4. \quad **if** \( \lambda^*(qr) \neq k/2 \) **then**
5. \quad \quad **insert** all leaves \( r \in R_i \) with edges \( qr \) and \( \lambda(qr) = \lambda^*(qr) \)
6. \quad **else**
7. \quad \quad **insert** a node \( q' \) and the edge \( qq' \) with \( \lambda(qq') = \lambda^*(qr) \)
8. \quad \quad **insert** all leaves \( r \in R_i \) with edges \( q'r \) and \( \lambda(q'r) = 0 \)
9. **end if**
10. **end for**

For an illustrative example see Fig. 3.

![Figure 3](image)

**Figure 3:** The graph \( G \) has two non-trivial false twin classes \( R_1 = \{2, 2'\} \) and \( R_2 = \{4, 4', 4''\} \). Its R-thin quotient \( G/\sim \) is explained w.r.t. \( \sim \) by the tree \((T^*, \lambda^*)\). The tree \((T_2, \lambda_2)\) is obtained from \((T^*, \lambda^*)\) using first alternative in Alg. 1 (line 5), amounting to replacing 2 by \(|R_1|\) leaves with the same parent. In the second step, vertex 4 is replaced by the subtree according to the second alternative in Alg. 1 (lines 7 and 8).

**Theorem 11.** \( G \) can be explained w.r.t. \( \sim \) if and only \( G/\sim \) can explained w.r.t. \( \sim \). If \((T^*, \lambda^*)\) is a canonical tree explaining \( G/\sim \), then a canonical tree \((T, \lambda)\) explaining \( G \) is obtained by Algorithm 1.
Lemma 12. There are unique labelings \( \lambda_{P_3} \) and \( \lambda_{K_3} \) of the tree \( S_3 \) with discrete \( \frac{3}{4} \) that explain the graphs \( P_3 \) and \( K_3 \), respectively:

\[
\lambda_{P_3}(e_1) = \lambda_{P_3}(e_2) = 2 \quad \text{and} \quad \lambda_{P_3}(e_3) = 0;
\lambda_{K_3}(e_1) = \lambda_{K_3}(e_2) = \lambda_{K_3}(e_3) = 1;
\]
Indeed explains the path obtained an analogous statement is true for \( \lambda \) non-connected graph \( \lambda \) and \( u \) argument implies the the remaining leaf \( \lambda \) order for \( y \) \( \lambda \). Proof. S \( \lambda \) by a 1-edge. The other neighbor of \( \lambda \) fully resolved tree on four vertices. Its interior edge explained by Lemma 14. The graph \( K_4+e \) can be explained by a fully resolved phylogenetic tree with a edge labeling explicitly given in the proof of Lemma 15. In contrast, \( C_4 \) and \( K_4-e \) cannot be explained with respect to \( \sim \) with an edge labeling with discrete \( 0 \) according to Lemma 15.

Proof. \( S_3 \) contains three paths on length two. Adopting the notation of Fig. 4 for both cases \( P_4 \) and \( K_3 \) we need \( \lambda(e_1) + \lambda(e_2) = 2 \) and \( \lambda(e_3) + \lambda(e_2) = 2 \). Therefore \( \lambda(e_1) \in \{0, 1, 2\} \). Explicitly enumerating the three cases yields:

- \( \lambda(e_1) = 0 \) implies \( \lambda(e_2) = 2 \) and thus \( \lambda(e_3) = 0 \), in which case \( x_1 \sim x_3 \), contradicting the fact \( \sim 0 \) is discrete.
- \( \lambda(e_1) = 1 \) implies \( \lambda(e_2) = \lambda(e_3) = 1 \), and thus \( G(\hat{2}) \) is \( K_3 \).
- \( \lambda(e_1) = 2 \) implies \( \lambda(e_2) = 0 \) and thus \( \lambda(e_3) = 2 \), whence \( G(\hat{2}) = P_3 \).

\( \square \)

**Lemma 13.** The path \( P_4 \) on four vertices \( x \sim y \sim z \sim u \) is explained only be the tree \( T = (xy)p \sim q(uz) \) with labels \( \lambda(xp) = \lambda(qu) = \lambda(pq) = 2 \) and \( \lambda(yp) = \lambda(qz) = 0 \).

Proof. First we observe that the path \( P_4 \) on four vertices cannot be explained by any labeling of a \( S_4 \). This leaves the fully resolved tree on four vertices. Its interior edge \( pq \) cannot be labeled 0. First consider \( \lambda(pq) = 1 \). It cannot contain an \( S_3 \) with all three edges labeled 1 since this would induce a triangle, i.e., at most one neighbor of \( p \), say \( x \), is attached by a 1-edge. The other neighbor of \( p \), call it \( y \), then must be attached by a 0-edge, since otherwise \( y \) is isolated. In order for \( y \) not to be isolated, \( q \) also must have a neighbor, that is connected via a 1-edge, say \( \lambda(qz) = 1 \). The same argument implies the the remaining leaf \( u \) must be connected to \( q \) with \( \lambda(qu) = 0 \). This tree, however, explains the non-connected graph \( K_2 \cup K_2 \). Thus \( \lambda(pq) = 2 \). Connectedness implies that at least one of the leaves attached to \( p \) and \( q \) must be labeled 0, say \( \lambda(py) = \lambda(qz) = 0 \), and thus \( y \) and \( z \) are adjacent in \( G \). It remains to consider the possible coloring for the remaining to edges \( \lambda(px) \) and \( \lambda(qu) \). If \( \lambda(px) = 1 \) then \( x \) is isolated for all choices of \( \lambda(qz) \). An analogous statement is true for \( \lambda(qy) = 1 \). If \( \lambda(px) = \lambda(qu) = 0 \) we obtain \( K_4-e \). If \( \lambda(px) = 0 \) and \( \lambda(qu) = 2 \) we obtain \( S_3 \). The same is true for \( \lambda(px) = 2 \) and \( \lambda(qu) = 0 \). Thus the only remaining choice is \( \lambda(px) = \lambda(qu) = 2 \). It indeed explains the path \( x \sim y \sim z \sim u \), see Fig. 5. \( \square \)

The fact that \( S_n \) is the only “exact-2-leaf root” of \( K_n \), i.e., the only tree with unit edge weights that explains \( K_n \) is shown in [10, Lemma 2]. It is not difficult to see that there is also no other choice of non-negative integer labels on \( S_n \) that explains \( K_n \).

**Lemma 14.** The complete graph \( K_n \) is explained with respect to \( \sim \) by the star \( S_n \) with the unique labeling function \( \lambda(e) = 1 \) for all \( e \in E(S_n) \).

Proof. Is is easy to check that this construction explains \( K_n \) for all \( n \geq 3 \). The trivial cases \( n = 1 \) and \( n = 2 \) are explained in the text. \( K_3 \) is only explained by \( S_3 \) with all edges labeled \( \lambda(e) = 1 \). Since the start \( S_n \) displays \( S_3 \) corresponding to every \( K_3 \) subgraph, all edges of \( S_n \) must be labeled by \( \lambda(e) = 1 \).

\( \square \)

We note for later reference that the uniqueness results in Lemmas 13 and 14 do not require the precondition that \( \sim \) is discrete. This observation will be important in the following section.
Lemma 15. There is no edge-labeled tree \((T, \lambda)\) with discrete \(\sim 0\) that explains the graphs \(C_4\) and \(K_4 - e\) with respect to \(\sim 2\). The graph \(K_3 + e\) is explained by a unique edge-labeled tree.

Proof. There are two topologically distinct trees for \(|L| = 4\), the star \(S_4\) and tree \(T_4\) with a single interior split. First consider the star \(S_4\). In order to explain \(K_3 + e\) or \(K_4 - e\) three of the four edges must be labeled 1 (corresponding to the induced \(K_3\). Depending on whether \(\lambda(e_1) = 1\) or \(\lambda(e_1) \neq 1\), the fourth vertex is either connected to all or none of the three other vertices. In order explain \(C_4\), there must be two edges with \(\lambda(e_1) = \lambda(e_2) = 2\) and one with \(\lambda(e_2) = 0\) corresponding to an induced \(P_3\). The remaining edge then must have \(\lambda(e_4) = 0\). But then \(x_2 \sim x_1\), contradicting that \(\sim 0\) is discrete.

Now consider the tree \(T_4\), which can be obtained from \(S_4\) by subdividing one of the edges and attaching an extra leaf to the subdividing vertex. Denote by \(s\) the (unique) inner edge of \(T_4\). Consider \(K_3 + e\) and \(K_4 - e\) as shown in Fig. 5. Then we must have \(\lambda(e_1) = \lambda(e_2) = 1\). If \(\lambda(s) = 0\) we recover the situation of \(S_4\), since the inner edge does not contribute to \(\sim 2\). On the other hand, if \(\lambda(s) = 2\), then \(x_1\) and \(x_2\) cannot be connected with \(x_3\) or \(x_4\), contradicting the existence of \(K_3\) as induced subgraph. Thus \(\lambda(s) = 1\). Then \(\lambda(e_3) = 0\). By assumption, \(\lambda(e_4) \neq 0\) since otherwise \(x_3 \sim x_4\). If \(\lambda(e_4) = 1\), the \(x_4\) is an isolated vertex in \(G\). If \(\lambda(e_4) = 2\), then \(x_4 \sim x_3\) while \(x_4\) is not in \(\sim 2\) relation to either \(x_1\) or \(x_2\). Thus \(G = K_3 + e\). The corresponding edge labeled tree is shown in Fig. 4. Since we have already considered all cases, \(K_4 - e\) cannot be explained with respect to \(\sim 2\).

Finally, consider \(T_4\) and suppose that \(G\) contains \(P_3\) as induced subgraph. There are two cases: If \(\lambda(e_1) = \lambda(e_2) = 2\) then connectedness of \(G\) implies that \(\lambda(s) = \lambda(e_3) = \lambda(e_4) = 0\), contradicting that \(\sim 0\) is discrete. In the alternative case we can assume, w.l.o.g., that \(\lambda(e_1) = 2\) and \(\lambda(e_2) = 2\). Furthermore, in order to explain \(C_4\) we must have \(\lambda(e_3) + \lambda(e_4) = 2\). If both \(\lambda(e_3) = \lambda(e_4) = 1\), then \(\lambda(s) = 0\) and \(\lambda(s) = 2\) yields \(G = K_2 \cup K_2\), for \(\lambda(s) = 1\) we obtain \(S_4\). In the remaining case we can choose \(\lambda(e_3) = 2\) and \(\lambda(e_4) = 0\). Now \(\lambda(s) = 0\) contradicts discreteness of \(\sim 0\), \(\lambda(s) = 1\) yields the edgeless graph. For \(\lambda(s) = 2\) we obtain \(P_4\). Thus \(C\) cannot be explained by \(T_4\) with respect to \(\sim 2\) by a labeling with discrete \(\sim 0\).

The fact that \(K_4 - e\) is a forbidden subgraph implies that two cliques in \(G\) cannot be “glued together” by a single common edges. It is possible, however, for cliques to touch in a cut vertex as shown by the example of the bowtie graph \(B\), which is obtained by gluing together two triangles at a common vertex, see Fig. 4.

Graphs that can be represented as pairwise compatibility graphs of caterpillars have received special attention in the literature [5, 33, 34, 35]. It is not difficult to see that the path \(P_h\), \(h \geq 3\) can represented by a caterpillar in several settings. These results cannot be directly applied in our setting, however. Any two leaves \(x\) and \(y\) attached to two distinct inner vertices of a caterpillar are separated by at least three edges an thus cannot be in relation \(\sim 2\) if we assume strictly positive integer weights. It follows immediately that \(P_h\) is not an exact-2-leaf power of a caterpillar and that \(P_h\) cannot be explained by caterpillar unless zero-weights are allowed. An explicit construction in [15] shows that \(P_h\) is explained by a caterpillar with edge weights in \([0, 1]\) with respect to exactly-1-relation \(\sim 2\). Lemma 4 implies that we can use the same construction to explain \(P_h\) by a caterpillar with edge weights in \([0, 2]\), see Fig. 6. It will be important later on that this construction is indeed unique:

Lemma 16. The path \(P_h\) has as its unique explaining tree the caterpillar \((T_h, \lambda_h)\) with all inner edges and the edges connecting to the end-points of \(P_h\) labeled 2 and all edges connecting to inner vertices of \(P_h\) labeled 0.

Proof. We first recall that the tree \((T_4, \lambda_4)\) explaining \(P_4\) is unique by Lemma 13. Now assume that for \(h \geq 5\), the tree \((T_{h-1}, \lambda_{h-1})\) explaining \(P_{h-1}\) is unique and thus a caterpillar. Any tree \((T, \lambda)\) explaining \(P_h\) therefore must display...
By Lemma 18,\( (T_h, \lambda_{h-1}) \), i.e., \((T, \lambda)\) is obtained from \(T_h\) by subdividing one edge and attaching leaf \(h\) and edge \(e_h\) to the new vertex, or by attaching \(h\) and \(e_h\) to an inner vertex of \(T_{h-1}\). One easily checks that the latter yields a branched tree or a disconnected graph. The same is true is any other edge except \(e_{p-1}\) and \(e_1\), the edges adjacent to the leaves \(h - 1 \) or 1 are subdivided. In the latter case, \(h\) cannot be adjacent to \(h - 1\). In the remaining case, the edge with which \(h - 1\) is attached is subdivided into an interior part \(s\) and the part \(e_{p-1}\) incident with \(h - 1\). Since the interior part cannot carry a zero label, we must have \(\lambda_h(e_{p-1}) = 0\), \(\lambda_h(s) = 2\), and \(\lambda_h(e_h) = 2\). Thus the caterpillar of Fig. 6 is indeed the only choice. We emphasize that this observation remains true even is \(0\) is not assumed to be discrete.

**Lemma 17.** The simple cycles \(C_p, p \geq 5\) cannot be explained with respect to \(\sim\) irrespective of whether \(0\) is discrete or not.

**Proof.** Every cycle \(C_p\) contains a path \(P_{p-1}\) with one vertex less as an induced subgraph. From Lemma 16 we known that \(P_{p-1}\) has a unique explanation by a caterpillar for all \(p \geq 5\). Thus any tree \((T^*, \lambda^*)\) explaining \(C_p\) thus must display caterpillar \((T_{p-1}, \lambda_{p-1})\) and thus \(T^*\) is obtained from \(T_{p-1}\) by either attaching \(p\) and \(e_p\) to inner vertex of \(T_{p-1}\) or by subdividing an edge and attaching \(p\) and \(e_p\) to the newly inserted vertex. As argued above, attachment to an inner vertex or subdivision of an edge other than \(e_1\) or \(e_{p-1}\) leads to a branched tree or a disconnected graph. If \(e_p\) is inserted by subdivision of \(e_1\), then \(p\) cannot be adjacent to \(p - 1\) and subdivision of \(e_{p-1}\) precludes adjacency of \(p\) and 1 for \(p \geq 3\). Thus the caterpillar tree \((T_{p-1}, \lambda_{p-1})\) cannot be extened to tree that explains \(C_p\) for any \(p \geq 4\). Note that this argument did not make the assumption that \(0\) is discrete.

Let us now turn to the general case. We first note that all graphs explained w.r.t. \(\sim\) with discrete \(0\) are chordal, i.e., every cycle of length greater than three has a chord. Even more stringently, every cycle of length 4 corresponds to a clique in \(G\) because the \(K_4 - e\), i.e., the 4-cycle with a chord, is also a forbidden induced subgraph. We note that there is ample literature on the relationship of chordal graphs and PCGs, see e.g. [4, 5]. Due to the differences in the edge weight functions, it is not immediately pertinent to our discussion, however.

**Lemma 18.** If \(G\) can be explained by the exact-2-relation with discrete \(0\) and contains a Hamiltonian cycle, then \(G\) is a complete graph.

**Proof.** The assertion is trivially true for \(n = 3\) and holds for \(n = 4\) because \(K_4\) and \(K_4 - e\), the only Hamiltonian graphs on 4 vertices except \(K_4\) are forbidden induced subgraphs. Now suppose the statement is true for for all \(|V| < p\) and consider a graph with \(p\) vertices. Since \(G\) is chordal, there is in particular a planar triangulation of \(C\) that is a subgraph of \(G\) and thus there are three consecutive vertices \(u - v - w\) along \(C\) such that \(u - w\) is also an edge in \(G\). Thus \(G\) is Hamiltonian. As an induced subgraph of \(G\) it can be explained by the exact-2-relation and thus is a complete graph by the induction hypothesis. Thus \(u - x - w\) is triangle in \(G\) and \(u - x - w - v\) is a cycle of length 4 in \(G\). Since \(C_4\) and \(K_4 - e\) cannot appear as induced subgraphs of \(G\), \([u, v, w, x]\) must for a clique in \(G\), and hence the edge \([v, x] \in E(G)\) for all \(x \in V(G)\). Thus \(G\) is a complete graph.

**Lemma 19.** A graph \(G\) with at least three vertices that can be explained by the exact-2-relation with discrete \(0\) is complete if and only if it is 2-connected.

**Proof.** If \(G\) is Hamiltonian, it is in particular also 2-connected. Now consider the case that \(G\) is 2-connected but not Hamiltonian. Let \(C\) be a cycle of maximal length in \(G\) and let \(x\) be a vertex not in \(C\). Then there is a cycle \(C'\) in \(G\) that contains \(x\) and at least two distinct vertices of \(C\) since otherwise one of the vertices of \(C\) would be a cut vertex of \(G\), contradicting 2-connectedness. Starting from \(x\), let \(p\) and \(q\) be first and last vertex of \(C\) encountered along \(C'\). By Lemma 18, \(G[C]\) is a complete graph, and hence there is a another Hamiltonian cycle \(C''\) in \(G[C]\) so that \(p\) and \(q\) are consecutive along \(C''\). Thus the cycle \(C''\) obtained traversing \(C''\) from \(p\) to \(q\) and then following \(C'\) from \(q\) through \(x\) back to \(p\) is a cycle that is strictly longer than \(C\), contradicting maximality. Thus \(G\) is Hamiltonian, and hence complete.

A graph \(G\) is a block graph [36] if each of its biconnected components is a clique. Lemma 19 thus implies that every graph that can be explained with respect to the exact-2-relations with discrete \(0\) is a block graph (see Thm. 21 below for a formal proof). Algorithm 2 (illustrated in Figure 7) explicitly constructs an edge-labeled tree that explains a given block graph.
Algorithm 2 Compute \((T(G), \lambda)\) for a connected block graph \(G\)

Require: a connected block graph \(G\)

1: mark “red” all cut vertices \(u \in V(G)\).
2: for all cliques \(K\) in \(G\) do
3: \(\text{if } K\text{ is an edge } e \text{ then}
4: \(\lambda(e) = 2\)
5: \(\text{else}
6: \text{replace } K\text{ by a star } S_{|V(K)|}\text{ with center } c_K\)
7: \(\lambda(u c_K) = 1\) for each \(u \in V(K)\)
8: \(\text{end if}
9: \text{end for}
10: \text{for all red vertices } v \text{ do}
11: \text{add a vertex } v'\text{ and edge } vv'\text{ with } \lambda(vv') = 0
12: \text{exchange the vertex names } v\text{ and } v'
13: \(\text{end for}
14: \text{return } (G, \lambda)\)

Lemma 20. Algorithm 2 transforms any connected block-graph \(G\) into an edge-labeled tree that explains \(G\) with respect to the exactly-2-relation \(\sim\) with discrete \(\bowtie\).

Proof. The output of Alg. 2 contains no cycles since all cycles in the input \(G\) are contained within a block and each block is replaced by a star. Furthermore, the replacement of a clique \(K_p\) by a star \(S_p\) with \(p + 1\) vertices preserves connectedness, hence \(G\) has been transformed into a tree at this stage. Every vertex of a clique \(K\), with \(|V(K)| \geq 3\) this is not also contained in another block is now a leaf; all other nodes of \(K\) are marked red. Every vertex in an \(K_2\) original \(K_2\) block is either a leaf or marked “red”. By construction, every “red” vertex has degree at least \(2\) and hence is not a leaf. The final operation adds a leaf to each “red” vertex. Together with the renaming of the vertices, thus, every vertex of the input graph is now a leaf in \(T\).

Now consider the labeling. First suppose that \(u\) and \(v\) are non-adjacent in the input \(G\), that is, there is a least one cut-vertex, say \(z\), between them in \(G\). The construction of \(T(G)\) ensures that the unique path from \(u\) to \(v\) in \(T(G)\) runs through a vertex \(z'\) that \(z\) as its neighbor. If the path from \(u\) to \(z\) in \(G\) ran through an edge in a triangle, it passes through the corresponding star and hence contains two edges labeled \(1\). Otherwise it runs through an unaltered \(K_2\)-block of \(G\), which is labeled \(2\). In each case, therefore, \(d_{T(G)\lambda}(x, y) \geq 4\). Now suppose that \(u\) and \(v\) are adjacent in \(G\). First suppose \(uv\) is contained in a triangle of \(G\). If neither \(u\) nor \(v\) was marked “red” they are both adjacent to the center \(c_K\) of a star with edges labeled \(1\). If \(u\) was a cut vertex, i.e., marked “red”, it appears a leaf adjacent to a vertex \(u'\) that in turn is adjacent to \(c_K\); furthermore \(\lambda(uu') = 0\) and \(\lambda(u c_K) = 1\). Analogous reasoning applied if \(v\) was a cut vertex of \(G\). In all cases, thus \(d_{T(G)\lambda}(uv) = d_{T(G)\lambda}(u, c_K) + d_{T(G)\lambda}(c_K, v) = 1 + 1 = 2\). If the edge \(uv\) is not contained in a triangle, then it is labeled \(2\). If \(u\) or \(v\) are cut vertices, then the unique path from \(u\) to \(v\) is \(u - u' - v\), \(u - v' - v\), or \(u - u' - v' - v\), with \(\lambda(uu') = \lambda(vv') = 0\) and a label \(2\) for the remaining edge. Hence, \(d_{T(G)\lambda}(uv) = 2\). In summary \(u \sim^2_{T(G)\lambda} v\) if and only \(u\) and \(v\) are adjacent in \(G\). Thus \(G\) is explained by \((T(G), \lambda)\) with respect to the exactly-2-relation.

Theorem 21. A graph \(G\) can be explained by an edge-labeled tree \((T, \lambda)\) with respect to the exact-2-relation with discrete \(\bowtie\) if and only if it is a block graph.

Proof. Suppose \(G\) can be explained w.r.t. \(R_t\) with discrete \(\bowtie\). If \(G\) is 2-connected, it is a clique by Lemma 19 and therefore also a block graph. Otherwise, we note that every 2-connected component \(G'\) of \(G\) is induced subgraph of \(G\) and thus, by Lemma 3, can be explained w.r.t. \(R_t\). By Lemma 19 every 2-connected component \(G'\) of \(G\) therefore must be a clique, i.e., \(G\) is a block graph.

Conversely, suppose that \(G\) is a block graph. Since Algorithm 2 is correct by Lemma 20, every connected block graph can be explained. Since the non-connected block graphs are just disjoint unions of connected block graphs, Lemma 5 completes the characterization of the non-connected case.
Figure 7: Illustration of Alg. 2. First all cut vertices of the input graph are marked. In the next step, edges not contained in triangles of $G$ receive label 2 and all larger cliques are replaced by stars with all new edges labeled 1. In the final step the original cut vertices are decorated by an additional neighbor with a 0-labeled edge.

The main result of this section is now obtained as

**Corollary 22.** A graph $G$ is explained by $\overset{2}{\sim}$ if and only if $R/\overset{2}{\sim}$ is a block graph.

*Proof.* Thm. 11 establishes that $G$ can be explained w.r.t. $\overset{2}{\sim}$ if and only if $R/\overset{2}{\sim}$ can be explained w.r.t. $\overset{2}{\sim}$. Since the graph $G/\overset{2}{\sim}$ is thin, Cor. 10 implies that $\overset{0}{\sim}$ is discrete for the canonical tree explaining $G/\overset{2}{\sim}$, and thus Thm 21 can be applied to $G/\overset{2}{\sim}$.

In the remainder of this section we consider the ambiguities in the construction of trees explaining block graphs. We start by characterizing contractible edges:

**Lemma 23.** Suppose $(T_\varepsilon, \lambda_\varepsilon)$ is obtained from a phylogenetic tree $(T, \lambda)$ by contracting the edge $e$ in $T$ and setting $\lambda_\varepsilon(e') = \lambda(e')$ for all $e' \neq e$ and suppose that $G(T, \lambda)$ is connected. Then $G(T, \lambda) = G(T_\varepsilon, \lambda_\varepsilon)$ if and only if $e$ is an interior edge of $T$ and $\lambda(e) = 0$.

*Proof.* We have already noted the contracting an inner 0-edge does not change the graph. By definition, leaf-edges cannot be contracted, since the vertices of $G$ correspond to the leaves of $T$. Connectedness of $G$ implies that there is a pair of vertices $x, y$ whose connecting path runs through $e$ and whose distance $d_{T, \lambda}(x, y) = 2$. The contraction of $e$ only leaves this distance unaffected if $\lambda(e) = 0$. Otherwise $d_{T_\varepsilon, \lambda}(x, y)$ changes, which implies that $x$ and $y$ become disconnected in $G'$ and hence the graph by the modified tree is different from $G$. 

Figure 8: Three alternative, topologically different canonical trees explaining the non-connected graph $2P_3$. These are derived from the unique tree explaining $P_3$ in Fig. 4.
We note that connectedness of $G$ is necessary in Lemma 23 since for non-connected $G$, the connected components can be “glued together” with arbitrarily complex trees as long as the distances between the attachment points is at least 3. In such examples it can be possible to contract edges without changing the explained graph. There are, for example at least three topologically different canonical trees that explain 2$P_3$, see Fig. 8.

Lemma 24. Let $(T, \lambda)$ be a canonical tree explaining a connected graph $G$ and let $x$ be an interior vertex in $T$. Then all edges incident to $x$ are 1-edges or $p$ has at least one adjacent leaf $u$ with $\lambda(\mu u) = 0$.

Proof. Suppose $p$ has no incident leaf. Since $G$ is connected, for every edge $e'$ there is another edge $e''$ such that $\lambda(e') + \lambda(e'') = 2$. For this pair of edges we have $\lambda(e') = \lambda(e'') = 1$ because no interior edges is 0-labeled. Thus $\lambda(e') = 1$ for all $e'$ incident with $x$. On the other hand, if $p$ has a neighbor $u$ with $\lambda(\mu u) = 2$, then connectedness of $G$ implies that there is another neighbor $y$ of $p$ with $\lambda(\mu y) = 0$. Hence, unless $p$ has only 1-neighbors, then there must be a least one incident 0-edge, which by assumption must be a leaf.

We remark, finally, that a tree with minimal number of vertices (or edges) that explains a graph with respect to the exactly-2-relation is necessarily canonical. Otherwise, the contraction of an edge would make it possible to decrease of both the number of edges and vertices.

4. Oriented Exactly-2-Relation

Generalizing the construction of the oriented exactly-1-relation in [15], we consider here an oriented version of the exactly-k-relation. In contrast to the previous sections, we consider here rooted trees $T$ with leaf set $L$. For $k$ leaves $x$ and $y$ there is a unique least common ancestor, denoted by $lca(x, y)$, defined as the vertex most distant from the root $r$ of $T$ that is common to the paths connecting $r$ with $x$ and $r$ with $y$, respectively.

Definition 25. Let $(\overline{T}, \lambda)$ be a rooted tree with leaf set $L$ and edge-labeling function $\lambda : E(\overline{T}) \rightarrow \mathbb{N}_0$. For $x, y \in L$ we consider the directed exactly-k-relation $\rightarrow_k$ defined by $x \rightarrow_k y$ if $\sum_{e \in \mathcal{P}(x, lca(x, y))} \lambda(e) = 0$ and $\sum_{e \in \mathcal{P}(lca(x, y), y)} \lambda(e) = k$ holds for the (unique) paths $\mathcal{P}(x, lca(x, y))$ from $x$ to $lca(x, y)$ and $(lca(x, y), y)$ from $lca(x, y)$ to $y$, respectively.

The rooted tree $(\overline{T}, \lambda)$ explains a the directed graph $\overline{G}(L, E)$ (with respect to the directed exactly-k-relation) if $(x, y) \in E(G)$ if and only if $x \rightarrow_k y$.

By construction $\overline{G}(L, E)$ is an oriented graph, i.e., at most one of $(x, y)$ and $(y, x)$ can be an edge. As in the unrooted case, we say that a rooted tree $(\overline{T}, \lambda)$ is canonical if it is a rooted phylogenetic tree and does not have an inner 0-edge. In the following we will consider the case that $\sim$ is discrete. As in the undirected case, we shall relax this requirement in the end.

As in [15], our strategy is to exploit the close relationships between the oriented and the undirected case. Therefore, we first derive some technical results regarding common properties of the oriented relation $\rightarrow_k$ and its undirected relative $\sim_k$.

Note that the underlying tree $(T, \lambda)$ of a rooted canonical tree $(\overline{T}, \lambda)$ is not necessarily an unrooted canonical tree. By contracting all the interior 0-edges and degree 2 vertices, we get a unique unrooted canonical tree $(T', \lambda')$ corresponds to $(\overline{T}, \lambda)$. Conversely, for any unrooted canonical tree $(T, \lambda)$ with $|V(T)| > 1$, we can create a set $\overline{T}(T, \lambda)$ of corresponding rooted least resolved trees as follows: (i) each interior vertex of $(T, \lambda)$ may serve as a root; (ii) each leaf attached by a 0-edge may serve as a root; and (iii) every 2-edge can be subdivided by inserting a root as a new vertex such that each of the two resulting edges is labeled 1. The construction is detailed in Algorithm 3. An example is given in Figure 9. The following lemma formalizes this one-to-one correspondence between unrooted canonical trees $(T, \lambda)$ and its corresponding sets of rooted canonical tree.

Lemma 26. Every rooted canonical tree can be constructed from its underlying unrooted canonical tree by Algorithm 3.
Algorithm 3 Compute the set of canonical rooted trees $(\mathbb{T}, \lambda)$

**Require:** unrooted canonical tree $(T, \lambda)$ with $|V(T)| > 1$

1: $(\mathbb{T}, \lambda) \leftarrow \emptyset$
2: **for** all interior vertices $v \in T$ **do**
3: designate $v$ as root
4: add the rooted tree to $(\mathbb{T}, \lambda)$
5: **end for**
6: **for** all leaf vertices $v \in T$ with $\lambda(vw) > 0$ where $N(v) = \{w\}$ **do**
7: subdivide $vw$ to $vw'w$ and designate $v'$ as root
8: relabel as $\lambda(vv') \leftarrow \lambda(vw)$ and $\lambda(v'w) \leftarrow 0$, designate $v'$
9: add the resulting rooted tree to $(\mathbb{T}, \lambda)$
10: **end for**
11: **for** all edges $e = uv$ with $\lambda(e) = k > 1$ **do**
12: subdivide the edge $e$ by inserting $v^*$ and designate $v^*$ as the root
13: **for** $j = 1..k-1$ **do**
14: $\lambda(vv^*) \leftarrow j$ and $\lambda(v^*v) \leftarrow k-j$
15: add the resulting rooted tree to $(\mathbb{T}, \lambda)$
16: **end for**
17: **end for**
18: **return** $(\mathbb{T}, \lambda)$

**Proof.** By construction, the set of canonical rooted trees corresponding to unrooted canonical tree is well defined, i.e., the correspondence is a mapping.

Suppose there are two distinct unrooted canonical trees $(T_1, \lambda_1)$ and $(T_2, \lambda_2)$ such that both their correspondings sets of rooted trees contains $(\mathbb{T}, \lambda)$. By construction, it has a underlying tree $(T, \lambda)$ from which a unique canonical tree is obtained by contracting 0-edges and degree 2 vertices. Thus $(T_1, \lambda_1) = (T_2, \lambda_2)$, a contradiction. Hence the mapping is injective.

The mapping is also surjective, since each rooted canonical tree $(\mathbb{T}, \lambda)$ can be constructed from its corresponding unrooted canonical tree.

**Lemma 27.** Suppose the unrooted canonical tree $(T, \lambda)$ explains $G$ with respect to $\mathbb{T}$. Let $\overrightarrow{G}$ be a digraph explained w.r.t. $\mathbb{T}$ by a rooted tree $(\mathbb{T}, \lambda)$ corresponding to $(T, \lambda)$. Then the underlying graph of $\overrightarrow{G}$ is a spanning subgraph of $G$.

**Proof.** By construction, $(T, \lambda)$ and $(\mathbb{T}, \lambda)$ has the same leaf set, and hence $V_G = V_{\overrightarrow{G}}$.

Any arc $x \rightarrow y$ in $\overrightarrow{G}$ is an edge in the underlying graph of $\overrightarrow{G}$ because that fact that $(\mathbb{T}, \lambda)$ explains $\overrightarrow{G}$ implies $\sum_{e \in P(x: lca(x,y))} \lambda(e) = 0$ and $\sum_{e \in P(y: lca(x,y))} \lambda(e) = k$. Considering the underlying unrooted graph $(T', \lambda')$, we have $\sum_{e \in P(x: y)} \lambda'(e) = k$. Since $(T', \lambda')$ explains $G$, and by construction $(T', \lambda')$ displays $(T, \lambda)$, we conclude that $(T', \lambda')$ also explains $G$. Hence $(x,y) \in E(G)$.

![Figure 9: Construction of rooted canonical trees A through F from the unrooted canonical tree on the left. The possible positions of the root are indicated by the triangles and the corresponding six rooted trees are shown to the right.](image-url)
Definition 28. Suppose there a tree \((T, \lambda)\) with discrete \(0\) that explains \(G\) w.r.t. \(\sim\). Then every subgraph \(H\) of \(G\) is allowed for \(\sim / 0\). Analogously, if there is a rooted tree \((\overrightarrow{T}, \lambda)\) with discrete \(0\) that explains \(\overrightarrow{G}\) w.r.t. \(\sim\), we say that every subgraph \(\overrightarrow{H}\) of \(\overrightarrow{G}\) is allowed for \(\sim / 0\).

In more detail a graph \(H\) is allowed for \(\sim / 0\) if there exists \((T, \lambda)\) such that for any \((x, y) \in E(G)\), we have \(\sum_{e \in P(x,y)} \lambda(e) = k\). If \(G\) is not allowed for \(\sim / 0\), we say that it is forbidden (as a subgraph) for \(\sim / 0\). Analogously, a graph \(\overrightarrow{H}\) is allowed for \(\sim / 0\) if there exists \((\overrightarrow{T}, \lambda)\) such that for any \((x, y) \in E(G)\), we have \(\sum_{e \in P(x,y)} \lambda(e) = k\) and \(\sum_{e \in P(u,v,x,y)} \lambda(e) = k\). If \(\overrightarrow{G}\) is not allowed as a subgraph in \(G\) w.r.t. \(\sim / 0\), we that say \(G\) is forbidden (as a subgraph) for \(\sim / 0\).

Lemma 29. If \(G\) is forbidden for \(\sim / 0\), then any orientation of \(G\) is forbidden in \(\sim / 0\). If \(\overrightarrow{G}\) is allowed as a subgraph for \(\sim / 0\) with rooted tree \((\overrightarrow{T}, \lambda)\), then its underlying graph is allowed for \(\sim / 0\) as a subgraph with the corresponding underlying tree \((T, \lambda)\).

Proof. Suppose, for contradictions, that \(G\) is forbidden for \(\sim / 0\) but the orientation \(\overrightarrow{G}\) of \(G\) is allowed for \(\sim / 0\). Then there exists a rooted tree \(\overrightarrow{T}, \lambda\) such that for any arc \(x \rightarrow y\) in \(G\), we have \(\sum_{e \in P(x,y)} \lambda(e) = 0\) and \(\sum_{e \in P(u,v,x,y)} \lambda(e) = k\) where \(u = lca(x, y)\). Consider the unrooted tree \((T, \lambda)\) of \((\overrightarrow{T}, \lambda)\). Since \((x, y) \in E(G)\) if and only if \(x \rightarrow y\) or \(y \rightarrow x\) is an arc in \(\overrightarrow{G}\), then for any \((x, y) \in E(G)\), \(\sum_{e \in P(x,y)} \lambda(e) = 0\) and \(\sum_{e \in P(u,v,x,y)} \lambda(e) = k\) where \(u = lca(x, y)\). By definition, \(G\) is allowed for \(\sim / 0\), i.e., we arrive at a contradiction. The second statement is a simple consequence of the first one.

The technical results obtained so far will allow us to infer properties of the oriented graph \(\overrightarrow{G}\) and their explaining trees \((\overrightarrow{T}, \lambda)\) from their underlying undirected graphs \(G\) and unrooted trees \((T, \lambda)\). In the following we will focus on graphs \(\overrightarrow{G}\) that can be explained w.r.t. \(\sim\) by a rooted tree \((\overrightarrow{T}, \lambda)\) with discrete \(0\).

Lemma 30. Oriented cycles are forbidden as a subgraph for \(\sim / 0\).

Proof. Suppose \(C_n\) is allowed. Then, by definition, there exists an orientation graph \(\overrightarrow{H}\) with vertex set \(V(C_n)\) such that \(\overrightarrow{C_n}\) is a subgraph of \(\overrightarrow{H}\) and a rooted tree \((\overrightarrow{T}, \lambda)\) that explains \(\overrightarrow{H}\). W.l.o.g., we assume \((\overrightarrow{T}, \lambda)\) is a rooted canonical tree.

Consider the underlying unrooted canonical tree \((T, \lambda)\) of \((\overrightarrow{T}, \lambda)\), we claim that \((T, \lambda)\) must be \((S_n, 1)\). Suppose that \((T, \lambda)\) explains graph \(G\). By Lemma 27 since the underlying graph \(H\) of \(\overrightarrow{H}\) is a subgraph of \(G\) which has the same vertex set with \(H\), and \(H\) contains a Hamiltonian cycle, thus \(G\) also contains a Hamiltonian cycle, by Lemma 18 \(G\) is a complete graph \(K_n\). By Lemma 14 the \((T, \lambda)\) displays \((S_n, 1)\).

Then we consider all the possibility to construct the set of rooted canonical trees corresponding to \((S_n, 1)\), and consider the oriented graph it explains. By Algorithm 3 we can place the root either on the center vertex, which will explain an empty graph, or place it on one of the leaves, which will explains oriented star on \(n\) vertices point to the leaves. In either case there is no cycles. By Lemma 26 we know we have constructed all rooted canonical trees and thus all oriented graphs they explain. Thus oriented cycles are forbidden as a subgraph for \(\sim / 0\).

Lemma 31. 2-star oriented to center; \(\bullet \rightarrow \bullet \leftarrow \bullet\), is forbidden as an induced subgraph for \(\sim / 0\).

Proof. Explicit construction shows that we obtain \(v_1 \sim v_2 \sim v_3\) for each of the three triples \(v_1v_2v_3, v_1v_3v_2, v_2v_3v_1\). This contradicts the assumption that \(0\) is discrete.

Lemma 32. Every graph \(\overrightarrow{G}\) that can be explained by an edge-labeled tree w.r.t. \(\sim / 0\) is an oriented forest with the property that all its component trees have a unique source vertex from which all arcs are directed away.
Proof. Let $\overrightarrow{G}$ be a graph that can be explained w.r.t. $\stackrel{2}{\rightarrow}/\stackrel{0}{\sim}$. Since all cycles are forbidden induced subgraphs, $\overrightarrow{G}$ is a forest. Furthermore, there is only a single source vertex in each connected component. Otherwise, if both $x$ and $y$ were sources within the same component tree, then the unique path from $x$ to $y$ would necessarily contain an induced subgraph of the form $\bullet \rightarrow \bullet \leftarrow \bullet$, which is forbidden.

A canonical tree $\overrightarrow{(T,\lambda)}$ with discrete $\stackrel{0}{\sim}$ that explains a connected oriented graph $\overrightarrow{G}$ w.r.t. $\stackrel{2}{\rightarrow}$ has a leaf that is attached to the root by 0-edge.

**Theorem 33.** $\overrightarrow{G}$ is explained w.r.t. $\stackrel{2}{\rightarrow}$ by a rooted tree $\overrightarrow{(T,\lambda)}$ with discrete $\stackrel{0}{\sim}$ if and only if $G$ is an oriented forest that does not contain the 2-star oriented to center, $\bullet \rightarrow \bullet \leftarrow \bullet$, as an induced subgraph.

**Proof.** To show the “only if” part, suppose $\overrightarrow{G}$ can be explained. Then $\overrightarrow{G}$ is oriented and by Lemma 30 it is an orientation tree, and by Lemma 31 we know that 2-star oriented to center is forbidden. For the “if” part we use the construction employed in [15] for $\stackrel{1}{\sim}$ (with 2-edges taking the place of 1-edges): To each inner vertex $v$ of $\overrightarrow{G}$ a new vertex $v'$ which represent $v$ in tree is attached with a 0-edge, while the inner edges of the tree have label 2. The Theorem now follows directly from Lemma 4.

We can relax the condition that $\overrightarrow{(T,\lambda)}$ has discrete $\stackrel{0}{\sim}$. To this end, we extend the false twin relation $x \preceq y$ to digraphs by setting $x \preceq y$ iff $x$ and $y$ have the same in- and out-neighbors. The quotient graph $G/\sim$ is known as the point-determining graph of $G$.

**Corollary 34.** An oriented graph $\overrightarrow{G}$ is explained w.r.t. $\stackrel{2}{\rightarrow}$ if and only if $\overrightarrow{G}$ is an oriented forest whose point-determining graph does not contain $\bullet \rightarrow \bullet \leftarrow \bullet$ as an induced subgraph.

**Proof.** It suffices to note that $\stackrel{0}{\sim}$-equivalent vertices are in the same $\preceq$-class and that there is a least resolved tree in which all members of a $\stackrel{0}{\sim}$-class are siblings.

We note, finally, that the rooted trees with discrete $\stackrel{0}{\sim}$ that explain $\overrightarrow{G}$ w.r.t. $\stackrel{2}{\rightarrow}$ are not unique, as exemplified in Figure 10.

### 5. Concluding Remarks

The main result of this contribution is the characterization of the exactly-2-relations, i.e. the graphs $\text{nniPCG}(T,\lambda,2,2)$. They form a proper superset of the the exact-2-leaf power graphs, which comprise only the disjoint unions of cliques. Section 2 suggests, however, that at least some of the structure and techniques carry over to general values of $k$. Several related problems are worth considering as well: in particular, $\text{nniPCG}(T,\lambda,k,\infty)$ and $\text{miniPCG}(T,\lambda,1,k)$ are of interest as models for coarse grained models of evolutionary distances.

The oriented version of the exactly-2-relation somewhat surprisingly, is much more closely related to the oriented exactly-1-relation of [15] that to the undirected exactly-2-relation. There is an alternative natural definition for a directed exactly-2-relation that omit the condition that $\sum_{e \in P(\lambda,\text{lca}(x,y))} \lambda(e) = 0$. Clearly, the resulting digraph are not
oriented, i.e., they may contain double edged. We suspect that their structure is more closely related to the Fitch graph (directed at-least-1-relation) recently studied in [17].

Regarding the analysis of rare-event data in phylogenetics the characterization of the exactly-2-relation naturally leads to the edge modification problem for block graphs and graphs whose R-thin quotient is a block graph, respectively. Although these problems do not seem to have been studied so far (see e.g. [37, Tab.1] and [38]). Since exactly-2-relation graphs are hereditary by Lemma 3, we suspect that edge modification problem for the exactly-2-relation graphs can be handled in manner similar to closely related edge modification problem for chordal graphs [37, 38] or cluster editing [39].

Acknowledgments

The authors gratefully acknowledge stimulating discussions with Marc Hellmuth, Manuela Geiß, and Maribel Hernández-Rosas on related classes of graphs derived from labeled trees. YJL acknowledges support of National Science Foundation for Young Scientists of China (Grant No.11901227).

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