Visibility Reasoning for Concurrent Snapshot Algorithms

Extended Version

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Visibility relations have been proposed by Henzinger et al. as an abstraction for proving linearizability of concurrent algorithms that obtains modular and reusable proofs. This is in contrast to the customary approach based on exhibiting the algorithm’s linearization points. In this paper we apply visibility relations to develop modular proofs for three elegant concurrent snapshot algorithms of Jayanti. The proofs are divided by signatures into components of increasing level of abstraction; the components at higher abstraction levels are shared, i.e., they apply to all three algorithms simultaneously. Importantly, the interface properties mathematically capture Jayanti’s original intuitions that have previously been given only informally.

1 INTRODUCTION

Linearizability [Herlihy and Wing 1990] is a standard correctness condition for concurrent data structures. It requires that the operations in any execution over the data structure may be ordered sequentially, without incurring changes to the observed results. In other words, the methods of a linearizable concurrent structure exhibit the same external behavior as the sequential equivalent. Programmers can use the concurrent variant to efficiently utilize modern systems’ multi-core setup, and rely on the sequential variant for understanding and formal reasoning.

Many methods exist for proving linearizability of a data structure. The standard idea shared by most of them involves finding, for each method of the structure, a point in time in the concurrent execution when the method can be considered as logically occurring. In other words, a method may execute over a period of time, admitting interference from other concurrent threads; nevertheless, for all reasoning intents and purposes, the execution is indistinguishable from one where the method executes atomically at a single point in time, without any interference. This point in time is referred to as the linearization point. Where the linearization point lies for a given method may depend on the run-time behavior and interleaving of other methods executing concurrently, sometimes even including behavior that occurs after the original method has already terminated (the latter are often termed far future linearization points). Although significant progress has been made recently in the verification of concurrent structures and algorithms, and in particular by the introduction of so-called prophecy variables [Abadi and Lamport 1991; Jacobs et al. 2018; Jung et al. 2020; Lynch and Vaandrager 1995] to model the dependence of linearization points on run-time behavior (including in the far future), establishing linearizability by explicitly exhibiting the linearization points of an algorithm remains a highly complex task in general.

A different approach, which avoids explicit reasoning about linearization points, has been proposed by Henzinger et al. [2013]. In this approach, one first specifies a set of properties and proves that the properties hold of every execution history over the data structure. One then constructs a proof of linearizability using only the specified properties as axioms, thereby abstracting from the underlying executions. In other words, the low-level reasoning about the concrete placement of linearization points is replaced by higher-level reasoning out of data structure axioms, in turn facilitating proof decomposition, modularity, abstraction, and reuse.

The properties used in the axiomatization are defined under relations tying an event in the execution to other events that depend on it, i.e., that observe it. For example, the event of reading of
a pointer will be related to the write event that was responsible for mutating the pointer. Once the pointer is mutated by another write, the subsequent reads will observe (i.e., be related to) the new write, or possibly another later write. Similar ideas of reasoning about event observations have been used in the axiomatizations of weak memory models [Raad et al. 2019] and in distributed systems [Viotti and Vukolić 2016] where they have been captured by means of so-called visibility relations. Reasoning by visibility has also been applied in different ways to prove linearizability of concrete data structures, such as queues [Henzinger et al. 2013] and stacks [Dodds et al. 2015], including reasoning automation [Bouajjani et al. 2017], and reasoning about relaxed notions of linearizability [Emmi and Enea 2019; Krishna et al. 2020].

In this paper, we demonstrate the further applicability of reasoning by visibility, by applying it in a novel way to even more nuanced algorithms and proofs. In particular, we show how visibility can be used to express and axiomatize the important internal properties shared by several snapshot algorithms. Similarly to the approach of Henzinger et al. [2013] to queues, the axiomatization enables a modularization of the linearizability proof: a significant portion of the proof is carried out once, and then reused for each snapshot algorithm.

More specifically, we verify the three snapshot algorithms by Jayanti [2005]. A snapshot algorithm scans a memory array and returns the values read, so that the obtained values reflect the state of the array at one point in time. In a sequential setting this is trivial to achieve since the array remains unchanged during the scan. However, Jayanti’s algorithms are concurrent, allowing interfering threads to modify the array while a scan is in progress. Jayanti’s algorithms are of increasing efficiency and generality. The simplest is the single-writer/single-scanner algorithm, which assumes that no two scanners run concurrently, and that no two writers concurrently modify the same array element. Jayanti’s second algorithm generalizes to a multi-writer/single-scanner setting, and the third is the most general and ultimately desirable multi-writer/multi-scanner version. Jayanti describes the linearization points only for the single-writer/single-scanner algorithm, but already this description is quite challenging to transform into a fully formal correctness proof [Delbianco et al. 2017] because the algorithm exhibits far future linearization points.

Of interest to us in the current paper is that each of the three algorithms builds on the previous one by relaxing some part of the previous algorithm’s implementation, while preserving the essential invariants that Jayanti calls forwarding principles. Jayanti credits the forwarding principles as the key idea behind his design, because the principles are shared by the three algorithms, and abstractly govern how a concurrent write into the array should be “forwarded” to a scanner that is in progress, but has already read past the written element.

In this paper, we show how the forwarding principles can be axiomatized mathematically in terms of visibility, which we present in Section 2 (Jayanti states the principles in English, much less formally). In Section 3, we develop the linearizability proof out of the axioms alone, so that it applies to all three algorithms simultaneously. We then establish that the axioms hold for each of the three algorithms: in Section 3 for the first algorithm and Section 4 for the second and third. In Section 5 we apply our method to verify another snapshot algorithm, that of Afek et al. [1993].

By employing visibility, we sidestep the difficulties inherent in reasoning about linearization points in general, and far future linearization points in particular. Our development substantiates that visibility is a natural abstraction to use in specifications and proofs, as it enables formally capturing the intuition—that of forwarding principles—that motivated the design of Jayanti’s algorithms in the first place. In summary, our contributions are as follows.

- This is the first formal proof of all of Jayanti’s three snapshot algorithms. Moreover, it efficiently reuses proofs between algorithms. Delbianco et al. [2017] employed a variation of the linearization point approach to prove only Jayanti’s first algorithm, with no direct way
of extending the proof to the other two algorithms, which are significantly more involved. That proof was mechanized in Coq. Based on loc. cit., Jacobs [2018] developed a mechanized proof of Jayanti’s first algorithm in VeriFast, using prophecy variables.

• We have axiomatized the forwarding principles, which were Jayanti’s motivating insight, and key properties of his snapshot algorithms, but have so far been out of reach of formal mathematics. The axiomatization is non-trivial, and required generalizing from Jayanti’s English description in order to apply to all three algorithms. It also enabled us to prove, for the first time, that forwarding principles formally imply linearizability. While visibility relations have been used before to axiomatize concurrent structures, this shows that they can also usefully capture important internal properties.

2 OVERVIEW

2.1 Jayanti’s First Snapshot Algorithm

The n-snapshot data-structure is an array of length n, which consists of two kinds of operations: scan which reads the memory and returns a list of length n reflecting the state of the memory at a point in time; and write(i, v) which writes the value v into memory cell i. We use s to range over instances of scan, w1 to range over instances of write(i, v) for some v, and e to range over all operations in general. In line with standard terminology of linearizability, we call these operations abstract events or abs events and color them magenta.

A snapshot (or any other) algorithm is linearizable if for every concurrent execution, there exists some way of sequentially ordering the overlapping (abs) events so that the value returned by each of the events remains unchanged compared to the concurrent execution, and moreover, matches the intended semantics of the data structure. Intuitively, linearizability implies that we can re-run the computation sequentially to obtain the same results as in the original concurrent run; however, the internal state of the algorithm during and after the runs need not match.

Implementing an efficient and correct concurrent snapshot algorithm is more challenging than it may seem. To highlight this point, consider a naïve implementation, where writers simply write to a shared array A and scanners simply iterate over the array A to obtain a snapshot. To allow this implementation to operate efficiently, we allow writes and scans to run concurrently. For this simple implementation, consider a snapshot array of length 2, where we have a scan s running concurrently with writes w0 writing 2 and w1 writing 3. Let 0 be the initial value of each array cell, and consider the following execution:

• Scan s starts and reads 0 from A[0], after which the scheduler interrupts s.
• Write w0 starts, writing 2 to A[0], and after w0 finishes, write w1 writes 3 to A[1].
• Scan s resumes, reading 3 from A[1]. It then returns the snapshot (0, 3).

The snapshot (0, 3) should indicate that there exists a point in time when the array consisted of that pair, however that is not the case. The array started as (0, 0), followed by (2, 0) after write w0, and (2, 3) after write w1, but none of these states are reflected in the result. In a sense, the scan s missed the write w0, yet it caught the write w1, which occurred after w0. Jayanti’s snapshot algorithms ensure that writes are not missed by the scanner, as we explain next.

Algorithm 1 is the first and simplest of Jayanti’s snapshot algorithms. The idea is for a scan to make two passes over the memory, first over the main array A, and then over the auxiliary array B. A writer updates the array B if it detects a concurrent scan via the boolean flag X, to forward its value. That is, in case the scanner missed the writer’s value when scanning A, it will have a chance to catch the value when scanning B. We thus refer to B as the forwarding array. Algorithm 1 is a single-writer/single-scanner algorithm, meaning that for it to behave correctly, two scans must not run concurrently and two writers must not concurrently mutate the same array cell. In
an implementation, this can be enforced by explicit locking, which we elide from Algorithm 1 following Jayanti’s original presentation.

Describing the algorithm in more detail, scan works by first setting $X$ to $true$ by event $s_{on}$, signaling that a scan is running the first pass. This is followed by clearing the forwarding array $B$ by setting all its cells to $\bot$ via the $sr_i$ events, for each $i$. The clearing ensures that the current scan cannot consider the forwards left over from the previous scans. Next, the scanner creates a naïve snapshot by copying the main array $A$ by the events $sa_i$ for each $i$ into the local array $V$. However, as we argued before, this is insufficient for a correct snapshot. This is where the second pass comes in, which starts after $X$ is set to $false$ by event $s_{off}$. In the second pass, the procedure repairs the naïve snapshot by stepping through the forwarding array $B$ by the events $sb_i$ for each $i$. If a non-$\bot$ (i.e., forwarded) value is found, it overwrites the original value in $V$, thus repairing the snapshot and preventing missing writes. Finally, the scanner returns $V$, which contains the complete snapshot.

For $write(i, v)$, it starts by writing its value to $A$ by event $wij_a$, followed by a check for a concurrently running scanner performing the naïve pass of the scan by event $wij_x$. If such a scan is detected, the writer forwards its value to $B$ by event $wij_b$.

Events marked in blue are called representation or rep events, and are used internally in the implementation of abs (i.e., magenta) events. The distinction between abs and rep events is standard in the theory of linearizability [Herlihy and Wing 1990]. We will use the naming convention whereby abs events and rep events with the same priming belong together, e.g. we assume $sa_i$ belongs to $s$ and $w’ij_a$ belongs to $w’j$.

Now, consider the scenario illustrated in Fig. 1, where we have Algorithm 1 being executed over an array of length two, with each value initially set to 0, with write events $w_{0j}$, $w’{0j}$ and $w_1j$ writing 2, 3 and 4 respectively, and a scan $s$. When $s$ reads $A[0]$ with $sa_0$ for its first pass, it reads 0, while when it reads $A[1]$ with $sa_1$ it reads 4 written by $w_1$ with $w_1a$. Between the two passes of $s$, $w_{0j}$ is missed, however since $X$ is set to $true$ by $s_{on}$, $w_{0j}$ will forward the value 2 with $w_{0j}b$. At the second
pass, \( s \) reads \( B[0] \) with \( s_b_0 \), reading 2 written by \( w_0 b \), and when it reads \( B[1] \) with \( s_b_1 \), it reads \( \perp \) written by \( sr_1 \), meaning \( s \) will use the original value 4 for its final snapshot, thus returning (2, 4).

In contrast to (0, 3) before, linearizability admits (2, 4) as a correct snapshot even though \( A \) never contained (2, 4) during the execution. Indeed, \( A \) only contained (0, 0), (2, 0), (3, 0) and (3, 4). This is actually fine, because we can reorder the concurrent events with the order \( w_0 \rightarrow w_1 \rightarrow s \rightarrow w' \), which, when executed sequentially, result in \( s \) having snapshot (2, 4). In the physical execution, the event \( w_0 \) returned before \( w'_0 \) and \( w_1 \) started. The reordering respects this by listing \( w_0 \) before \( w'_0 \) and \( w_1 \). In other words, the reordering affects only events that physically overlapped, as required by linearizability.

Jayanti sketches the correctness proof of Algorithm 1 by describing its linearization points. The linearization point for a scan \( s \) is always when the scan performs \( s_{off} \). However, the linearization point of a write \( w_i \) varies. If there is no scan concurrent to \( w_i \), or there is a concurrent scan, but it reads the value of \( w_i \) either from \( A \) or from \( B \), then the writer’s linearization point is at \( w_i a \). If there is a concurrent scan \( s \) that misses \( w_i \), which can occur if the scanner misses \( w_i \) in its \( A \) pass, and \( w_i \) either does not write into \( B[i] \) due to \( X \) being set to \( false \) before the writer could forward, or \( w_i \) writing into \( B[i] \) too late, then \( w_i \) must be logically considered as occurring after \( s \). Thus, the linearization point of \( w_i \) is immediately after \( s_{off} \), making the linearization point of write \( external \), as its position is given in terms of another procedure, in this case scan. The observation that a scan missed a write \( w_i \), which occurs when the scan reads \( B[i] \) with \( s_b_i \), can be made after the writer has already terminated, making write \( external \) exhibit a far future linearization point.

We proceed to show how to organize the linearizability proof of Algorithm 1, and the other two Jayanti algorithms, by axiomatizing forwarding via visibility, without using linearization points.

### 2.2 Basic Abstractions of Visibility Reasoning

#### 2.2.1 Events and Their Structure

An event is an object consisting of fields start, end, op, in, and out, describing the following aspects of the execution of some operation of the data structure: start is the event’s beginning time, end is the ending time, op is the operation name (e.g., scan or write), in is the operation’s input, and out is the output. We refer to the elements of an event \( e \) by projection, e.g. \( e.start \) and \( e.op \). The fields \( e.in \) and \( e.out \) are natural numbers, and \( e.end \) may be \( \infty \) (infinity) to represent that \( e \) has not terminated yet, i.e., \( e \) is terminated iff \( e.end \neq \infty \), which we denote with \( \mathcal{T}(e) \). For every \( e \), \( e.end \geq e.start \). The types of \( e.in \) and \( e.out \) depend on \( e.op \), and \( e.out \) is undefined iff \( e.end = \infty \).

Additionally, each event \( e \) contains the optional field parent, corresponding to the event that invoked \( e \), if any. For example, if \( e \) is a \( rep \) event, then \( e.parent \) is the abs event that contains \( e \). Note that \( e.parent \) and \( e \) need not have the same \( op \), in and out fields; e.g., a write \( rep \) event can be invoked both by abs writer and abs scanner.

Finally, we require that each abs event is single-threaded, and thus cannot fork children threads.

We denote the set of all events of a given execution history by Events. If \( E \subseteq \text{Events} \), then \( op(E) = \{ e \in E \mid e.op = op \} \) selects the events with operation \( op \), and we overload \( \mathcal{T} \) over sets with \( \mathcal{T}(E) = \{ e \in E \mid \mathcal{T}(e) \} \) to select the terminated events.

As linearization order can only affect physically overlapping events, proving linearizability requires reasoning about non-overlapping events, which is captured by the \( \text{returns-before} \) relation

\[
e \sqsubseteq e' \iff e.end < e'.start
\]

denoting that \( e \) terminated before \( e' \) started. Two events are overlapping if they are unrelated by \( \sqsubseteq \). As customary, we write \( \sqsubseteq \) for the reflexive closure of \( \sqsubseteq \). The relation \( \sqsubseteq \) is irreflexive (i.e. acyclic)
We write \( e \triangleq \) happens-before we define the \( \prec \) partial order, and moreover, an interval order [Felsner 1992], as it satisfies the following property

\[
e_1 \sqsubset e_2 \land e_1' \sqsubset e_2' \implies e_1 \sqsubset e_2' \lor e_1' \sqsubset e_2
\]  

(RB.1)

Fig. 2 illustrates why Property (RB.1) must hold. The figure shows events \( e_1, e_2, e_1', e_2' \) such that \( e_1 \sqsubset e_2 \), and \( e_1' \sqsubset e_2' \), and Property (RB.1) holds because also \( e_1 \sqsubset e_2' \). We can try to invalidate the latter by shifting \( e_1' \) and \( e_2' \) to the left so that \( e_2'.start < e_1.end \) while maintaining \( e_1'.end < e_2'.start \). But then we’re forced to have \( e_1'.end < e_2.start \), i.e. \( e_1' \sqsubset e_2 \) which re-establishes Property (RB.1).

We also say that \( e \) is a subevent of \( e' \) (alternatively, \( e' \) contains \( e \)) if

\[
e \sqsubseteq e' \iff e'.start \leq e.start \land e.end \leq e'.end
\]

For any rep event \( e \), we require that \( e.parent \) must be an abs event such that \( e \sqsubseteq e'.parent \). Additionally, the following property holds for subevents and returns-before.

\[
e_1 \sqsubseteq e_1' \land e_2 \sqsubseteq e_2' \land e_1' \sqsubset e_2' \implies e_1 \sqsubset e_2
\]  

(RB.2)

That is, if \( e_1' \) returned before \( e_2' \) then all subevents of \( e_1' \) must return before any subevent of \( e_2' \). For example, in Fig. 2, \( e_1 \sqsubseteq e_1' \) and \( e_2 \sqsubseteq e_2 \), and \( e_1 \sqsubset e_2 \). If we shift \( e_1' \) to the left so that \( e_1' \sqsubset e_2 \), then we just increase the distance between \( e_1 \) and \( e_2' \), maintaining \( e_1 \sqsubset e_2 \).

2.2.2 Visibility Relations. If the event \( e' \) depends on the result of \( e \), we say that \( e \) is visible to \( e' \), or alternatively that \( e \) is observed by \( e' \). We denote the relationship as

\[
e < e'
\]

Depending on the data structure being verified, we will often require several different observation relations to differentiate how the observation came about. For example, in the case of Jayanti, we will use \( \text{rdw}_f \) for a “reads-from observation” (a reader observes a writer by reading what was written), and \( \text{fwd}_w \) for a “forwarding observation” (a scanner observes a writer by having the written value forwarded). The former relation is tied to the physical act of reading a pointer, and will be the same in all three algorithms. The latter relation differs for different algorithms and is typically provided by the human verifier, similarly to how loop invariants must often be provided.

We will typically obtain the visibility relation \( < \) by unioning all the different observation subrelations. We then require the following property of \( < \)

\[
e < e' \implies e' \not\sqsubseteq e
\]  

(V.1)

where \( <^+ \) is the transitive closure of \( < \). Read contrapositively, the property says that if \( e' \) terminated before \( e \) started, then \( e' \) cannot end a non-empty sequence of observations starting from \( e \). In particular, as a special case, \( e' \) cannot observe \( e \) or be equal to \( e \), meaning \( < \) has to be irreflexive. We write \( \prec \) and \( \sqsubseteq_f \) for the reflexive closure of \( < \) and \( \sqsubseteq \) respectively.

2.2.3 Happens-Before Relation. Given two events \( e \) and \( e' \), if \( e \sqsubset e' \) or \( e < e' \), then clearly, in the ultimate linearization order we want to construct, \( e \) must appear before \( e' \). To capture this intuition, we define the happens-before relation as the transitive closure

\[
\prec \triangleq (\sqsubset \cup <)^+
\]
We also name the single-step happens-before relation \( \ll \subseteq (W \cup R)^2 \) as \( \ll \), where \( \ll \) is the reflexive-transitive closure of \( \ll \). The \( \ll \), \( \ll \), and \( \ll \) relations are all standard in the literature [Viotti and Vukolić 2016]. An important property is that \( \ll \) is an irreflexive (i.e., acyclic) partial order, which is ensured by \( \ll \) satisfying Property (V.1).  

**Lemma 2.1.** If visibility relation \( \ll \) satisfies Property (V.1) then \( \ll \) is irreflexive.

**Proof.** We assume that \( \ll \) is not irreflexive (i.e., there exists \( e \) such that \( e \ll e \)), and derive contradiction. The relation \( e \ll e \) is a cyclic chain of \( \ll_1 \), each of which is either \( \ll \) or \( \ll \). We are justified in considering chains with only one or zero occurrences of \( \ll \), as chains with more occurrences of \( \ll \) can be shortened, thus we can recursively shorten a chain until it consists of one or zero \( \ll \). Indeed, if the chain has more than one \( \ll \), it has the form

\[
e \ll_1 \cdots \ll_1 e_1 \ll e_2 \ll \cdots \ll e_3 \ll e_4 \ll_1 \cdots \ll_1 e,
\]

where \( e_2 \) and \( e_3 \) are related by a chain of zero or more \( \ll \)'s (i.e., \( e_2 = e_3 \) or \( e_2 \ll^* e_3 \)). But then, we can remove one \( \ll \) as follows. If \( e_2 = e_3 \), by transitivity of \( \ll \), we can shorten the chain to \( e \ll e_1 \ll e_4 \ll e \). If \( e_2 \ll^* e_3 \), by Property (RB.1), it must be either \( e_1 \ll e_4 \) or \( e_3 \ll e_2 \). If \( e_1 \ll e_4 \), we again shorten to \( e \ll e_1 \ll e_4 \ll e \). Otherwise, \( e_3 \ll e_2 \) and \( e_2 \ll^* e_3 \) contradict Property (V.1).

On the other hand, if there is exactly one \( \ll \) in the chain, i.e., \( e \ll^* e_1 \ll e_2 \ll^* e \), then \( e_2 \ll^* e \ll^* e_1 \), i.e., \( e_2 \ll^* e_1 \) which, along with \( e_1 \ll e_2 \), contradicts Property (V.1). Finally, if there is no \( \ll \) in the chain, i.e., \( e \ll^* e \), then Property (V.1) directly implies the contradiction \( e \ll^* e \).

Fig. 3. Signature representation of the properties of an atomic register.

We will usually color the relations with the same colors as the events they are relating (e.g., \( e \ll e' \) for abs events \( e \) and \( e' \), and \( e_r \ll e_r' \) for rep events \( e_r \) and \( e_r' \)).

### 2.2.4 Memory Model

With the above relations we can now state as axioms the properties that we expect of the underlying memory model. We start with a simple axiomatization of memory that is sufficient for Algorithm 1. We will extend this axiomatization in Section 4 to account for the LL (load-link) and SC (store-conditional) memory operations, required by Algorithms 2 and 3.

The memory model only considers operations over individual memory cells, aka. atomic registers [Herlihy and Shavit 2008]. Their axiomatization in terms of visibility relations is given in Fig. 3, in the form of a signature we assume the memory to satisfy. We use blue in this figure, since we will only use events of atomic registers as rep events. The sets \( W \) and \( R \) are defined by the
hypothesis as the sets of writes and reads respectively. The quantifier $\Sigma$ signifies that the relations $\prec$ and $\sqsubseteq$, are abstract components of the specification; the clients do not know anything about these components outside of the listed axioms. Given an instance $X$ of the signature, the clients can refer to the relations and to the let definitions by projection as $X.\prec$, $X.\sqsubseteq$, $X.W$ and $X.R$. We will present similar signatures for snapshot data-structures and Jayanti-style forwarding.

Looking at Fig. 3, an atomic register is mathematically represented by two sets of events: $W$, corresponding to all the writes into the register’s memory cell, and $R$, corresponding to all the reads. These are related by $w \sqsubseteq r$, stating that read $r$ reads from (i.e., observes) write $w$. These events may also be related by some other relation part of the total set of observations $\prec$, which includes $\sqsubseteq$; Property (V.1) must hold of $\prec$ to ensure that $\prec$ is irreflexive (thus, a partial order) as per Lemma 2.1. Property (M.1) states that for any terminated read $r$ there exists an observed write $w$ with its input being the same as the output of $r$. Property (M.2) states that the read $r$ can observe only the latest write $w$ into the given memory cell. Property (M.3) states that a read may observe at most one write. Property (M.4) states that the writes into the given cell are totally ordered. These are standard properties of sequentially consistent memory [Herlihy and Shavit 2008].

We want to reason about $\prec$ order arbitrarily between rep events, even if they belong to distinct register objects. As an analogue to the locality property of linearizability, which says that the union of multiple linearizable objects is itself linearizable, we establish that the union of objects satisfying Property (V.1) itself satisfies Property (V.1). By Lemma 2.1, this implies that $\prec$ order between any rep event is a partial order. More formally, let $R_1 \ldots R_m$ be all objects satisfying $\text{AREg}$ and representing our memory. We define the top-level $\prec$ (and top-level $\sqsubseteq$) relation over rep events as the union of visibility of all register objects:

\[
\prec \triangleq (\sqsubseteq \cup <) \cup (R_1.\prec) \cup \cdots \cup (R_m.\prec)
\]

Since each register satisfies Property (V.1) and each individual $\prec$ only relates events from the same register, it follows that the combined $\prec$ also satisfies Property (V.1). By Lemma 2.1, the top-level $\sqsubseteq$ is then a partial order. The combination similarly satisfies Properties (M.1) to (M.4). We will use these top-level definition when relating rep events of distinct objects.

### 2.2.5 Linearizability.

We now define linearizability formally in terms of visibility relations.

**Definition 2.2.** History $E$ is linearizable with respect to data structure $D$ if there exists a visibility relation $\prec$ and a total order $\prec$ on the set of events $E_c = \{e \mid \exists e' \in T\ (E), e < e'\}$, such that: (1) $\sqsubseteq = (\sqsubseteq \cup <)^*$ is a total order on $E_c$, and (2) executing the events in $E_c$ in the order of $\sqsubseteq$ is a legal sequential behavior of $D$. The order $\prec$ is the linearization order (or linearization, for short) of $E$. Algorithm (or structure) $A$ is linearizable wrt. $D$ if every history of $A$ is linearizable wrt. $D$.

Definition 2.2 differs somewhat from the original one [Herlihy and Wing 1990] in that we use the visibility relation $\prec$ to complete $T\ (E)$ into $E_c$ with events that have been observed, but have not yet terminated. The original definition existentially abstracts over the set of completing events; formally, it does not organize them into a visibility relation, but in practice these events are always added because they have been observed by some terminated event, and are necessary to ensure legal sequential behavior. The requirement $\sqsubseteq \subseteq <$ means that $\prec$ must respect $\sqsubseteq$ (in addition to $\prec$), and in particular that $\prec$ can reorder only overlapping events. In this paper, $D$ is the snapshot data structure that has the following sequential behavior over its state (the array $A$):

- $w_i \in \text{WRITE}_i$ writes $w_i.\text{in}$ to $A[i]$ and returns nothing, i.e., $w_i.\text{out} = ()$.
- $s \in \text{SCAN}$ does not modify $A$ and returns a copy of $A$, i.e., $s.\text{out} = A$. 
2.3 Hierarchical Structure of the Proof

Ultimately, our goal is to prove that each of Jayanti’s three snapshot algorithms is linearizable. To accomplish this in a way that maximizes proof reuse, we divide our proofs into multiple segments, which we illustrate in Fig. 4. The proof of Jayanti’s first algorithm is split into the following steps:

(1) Jayanti’s single-writer/single-scanner algorithm (Algorithm 1) consists of the rep events described in Fig. 6 and satisfies the forwarding snapshot signature that we give in Fig. 7.

(2) Any algorithm consisting of the rep events described in Fig. 6 and satisfying the forwarding snapshot signature from Fig. 7, also satisfies the general snapshot signature in Fig. 5.

(3) Any algorithms satisfying the general snapshot signature in Fig. 5 is linearizable.

For Jayanti’s remaining algorithms, we incorporate the following two additional steps.

(1a) Jayanti’s two multi-writer algorithms (Algorithms 2 and 3) consist of the rep events described in Fig. 10 and satisfy the multi-writer forwarding snapshot signature in Fig. 11.

(2a) Any algorithm consisting of the rep events described in Fig. 10 and satisfying the multi-writer forwarding snapshot signature in Fig. 11, also satisfies the forwarding snapshot signature in Fig. 7. The linearity then follows by steps (2) and (3) above, which are thus reused and shared by all three algorithms.

We proceed to describe the intuition behind the steps (1) to (3), together with the forwarding and snapshot signatures. The detailed proofs of (1) to (3) are in Section 3. The description and proofs of (1a) and (2a) are in Section 4.

2.4 Snapshot Signature

Going backwards, we start with the general snapshot signature (Fig. 5). It describes the axioms for reasoning about the snapshot data structure as a whole, and will interface the linearity proofs by the following lemma, which we will prove in Section 3.

**Lemma 2.3.** Histories satisfying the Snapshot signature (Fig. 5) are linearizable.

The Snapshot signature is almost identical to that for atomic registers (Fig. 3); the main distinction is that we have multiple sets of writes into multiple pointers, scans (observing multiple writes) instead of reads, the sets of effectful writes $\mathcal{W}_i$, and the extra Properties (S.5) and (S.6).

Not all writes are immediately observable once they have started, therefore we introduce the set of effectful writes $\mathcal{W}_i$ for each $i$. These are the writes that are available for scans to observe (by definition of $\mathcal{W}_i$) and which we can order by Property (S.4). Property (S.5) encodes that each terminated write must have been effectful.

Property (S.6) imposes a form of monotonicity on the ordering of the writes observed by scans over multiple memory cells. More specifically, if we have two writes $w_i, w_j$ observed by scan $s$, and two writes $w'_i, w'_j$ observed by scan $s'$, then the writes into $i$ and $j$ cannot be ordered in the
signature | Snapshot ≜ Snapshot data-structure
---|---
Let | \( W_i \) Set of all writes of cell \( i \) of the snapshot in history
\( S \) Set of all scans of the snapshot in history
\( \sum \forall l \subseteq W_i \) Set of effectful writes of index \( i \)
\( < \subseteq (\bigcup_l W_i \cup S)^2 \) Visibility relation
\( \triangleright \subseteq < \cap \bigcup_l W_i \times S \) Reads-from visibility
Let \( \triangleright \triangleq (\sqcup \cup <)^+ \) Happens-before order
\( \forall e, e' \). \( e < e' \implies e' \not< e \) (V.1)
\( \forall w_i, s \in \mathcal{T}(S) \). \( \exists w_i, w_i \triangleright s \land w_i.in = s.out[i] \) (S.1)
\( \forall w_i, w_i', s. \quad w_i \triangleright s \implies \# w_i', w_i < w_i' < s \) (S.2)
\( \forall w_i, w_i', s. \quad w_i \triangleright s \land w_i' \triangleright s \implies w_i = w_i' \) (S.3)
\( \forall w_i, w_i' \in W_i. \quad w_i \neq w_i' \implies w_i < w_i' \lor w_i' < w_i \) (S.4)
\( \mathcal{T}(W_i) \subseteq W_i \) (S.5)
\( \forall w_i, w_i', w_j, w_j', s, s'. \quad w_i, w_j \triangleright s \land w_i', w_j' \triangleright s' \land w_i < w_i' \implies w_j' \not< w_j \) (S.6)

Fig. 5. Signature representation of the properties of a snapshot data-structure.

opposite way. The property ensures that the scans can be totally ordered in an eventual linearization order. Indeed, if \( w_i < w_i' \) and \( w_i' < w_j \), and we ordered \( s \) before \( s' \), then we reach a contradiction by Property (S.2), because the event \( w_j \) occurs between \( w_i' \) and the scan \( s' \) that observes it. Similar argument applies if we try to order the scans the other way around.

2.5 Forwarding Signature

The Snapshot signature in Fig. 5 captures the key properties of snapshot algorithms, and we shall prove linearizability solely out of the axioms of this signature. However, Jayanti’s algorithms have more in common than merely being snapshot algorithms, as they share the same design principle of forwarding. We can thus further modularize the proofs by axiomatizing forwarding itself.

To see what the axiomatization should accomplish, we must first discuss the high-level differences between Jayanti’s algorithms that the axiomatization must abstract over. We present Algorithms 2 and 3 in detail in Section 4, but for now it suffices to know that in order to support the multi-writer functionality of Algorithms 2 and 3, we need to allow for different kinds of forwarding that go beyond simply writing into the auxiliary array \( B \) that Algorithm 1 does.

Additionally, to support multi-scanner functionality of Algorithm 3, we need to allow for an abstract notion of *virtual* scan. While in a multi-scanner setting several physical scans may run concurrently, the key idea of Jayanti for Algorithm 3 is that these scans collaborate to create a virtual scan (and a corresponding snapshot), of which at most one may exist at any given moment. Thus, even though Algorithm 3 is physically a multi-scanner, in Section 4.3 we will still be able to conceptually see it as a single-scanner one.

While rep and abs events originate from execution histories, the virtual scans are artificial events that the human verifier creates themselves for purposes of verification. In analogy with the concept of “ghost state” that is frequently used in verification of concurrent programs, we can say that virtual scans are “ghost” events. We will further require a mapping that sends each abs scan to a virtual scan to which the abs scan contributed. Thus, in a multi-scanner algorithm, a number of abs scans may be mapped to the same virtual scan. The mapping may be partial because some abs scan
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Table:

| resource | A : array[n] of AReg | Main memory array |
|----------|---------------------|------------------|
| event signature | WRITE_i ≜ a : A[i].W | Rep event of writer writing value into A[i] |
| event signature | VSCAN ≜ B : array[n] of AReg | Forwarding array of this virtual scan |
| | r_i : B[i].W | Rep event of scanner resetting B[i] by writing ⊥ |
| | a_i : A[i].R | Rep event of scanner reading from A[i] |
| | b_i : B[i].R | Rep event of scanner reading from B[i] |

Fig. 6. Resources and event signatures for abs writes and virtual scans corresponding to Forwarding signature.

may not have yet reached a point for which the virtual scan representing it has been determined. For single-scanner algorithms, this map is a (total) bijection that identifies virtual and abs scans. In particular, the intuition about abs scans from Algorithm 1 will suffice to understand virtual scans in our axiomatization of forwarding principles. In the rest of the text, we use the color green to visually mark the virtual scans.

2.5.1 Event Signatures. For the Forwarding signature, it is not enough to just consider properties over the abs events, as for the Snapshot signature, but we also need to consider their rep events. For this reason, we introduce a special kind of signature, called event signature, to encode which rep events an abs or virtual event consists of. In particular, for an abs event e, we write e ∈ sig, if e’s rep events are enumerated by the fields of the signature sig. If e ∈ sig, where sig contains the field a, we write ea for the projection e.a. Similarly, if e is a virtual event. We shall further assume the following properties of structures satisfying event signatures:

(ES.1) Let e be an event in a signature with field a. If e is an abs event, then ea.parent = e, and thus also ea ⊆ e, by our assumption on event structure in Section 2.2.1. If e is a virtual event, we do not insist on ea.parent = e. Virtual event e is a custom collection of rep events, whose parents remain the abs events that invoked them, not the collection that is e. Nevertheless, we still require ea ⊆ e.

(ES.2) Any event in an event signature instance is unique to the instance, i.e., if e, e’ ∈ sig, and sig contains the field a, and ea = e’a, then e = e’.

(ES.3) Instances of event signatures may have some (or all) of their subevents undefined, unless the instance is a terminated event, in which case all its subevents must be defined. For example, if e is non-terminated, ea may be undefined. This corresponds to ongoing abs events having only a subset of its rep events executed so far.

We now present Fig. 6 which consists of the resources and event signatures necessary for the Forwarding signature. First, we have the array A which corresponds to the main memory that a forwarding snapshot algorithm operates over. Algorithm 1 clearly has such an array. Next, we have that every abs write w_i has a unique rep event w_i.a signifying the physical act of writing into A. This is clearly true of Algorithm 1, as we have seen. For a virtual scan σ, the figure postulates an array B, as well as three rep events σ_r_i, σ_a_i, and σ_b_i, corresponding to resetting B[i], reading from A[i], and reading from B[i], respectively. In Algorithm 1, we have already described these events under the names sr_i, sa_i, and sb_i, and similar events will exist for Algorithms 2 and 3 as well. The notation implies that the array B and the subevents are local fields of σ. We will make use of the locality of
| **signature** | **Forwarding** | **Snapshot data-structure with forwarding** |
|--------------|---------------|-----------------------------------------|
| **Let**      |               |                                         |
| \( W_i \)   | Set of all writes of cell \( i \) of the snapshot in history with each \( w_i \in \text{WRITE}_i \) | **Set of all scans of the snapshot in history** |
| \( S \)     |               |                                         |
| \( \Sigma \) | **Set of virtual scans with each \( \sigma \in \text{VSCAN} \)** | **Partial mapping of abs scans into virtual scans** |
| \([-\]^* \) | \( \rightarrow \) \( \Sigma \) | **Forwarding visibility** |
| \( \text{fwd}_{\rightarrow} \subseteq \bigcup_j W_j \times \Sigma \) | **Reads-from visibility** |
| **Let**      |               | **Writing visibility** |
| \( w_j \xrightarrow{r_j} \sigma \) | \( \triangleq \) | **Visibility relation** |
| \( w_i \xrightarrow{w_i} w_i' \) | \( \triangleq \) | **Happens-before order** |
| \( < \triangleq \) | \( \cup \) |                                         |
| \( \sim \triangleq \) | \( x \) |                                         |

\( \forall s \in \text{dom}([-\]^*). \) \( \triangledown s \leq s \) \( \text{(F.1)} \)

\( \forall i, s \in \mathcal{T}(S). \) \( \exists w_i, w_i' \xrightarrow{r_i} s^* \land w_i.\text{in} = s.\text{out}[i] \) \( \text{(F.2)} \)

\( \forall e_r \in A[i].W. \) \( \exists w_i, w_i.a = e_r \) \( \text{(F.3a)} \)

\( \forall \sigma, e_r \in \sigma.B[i].W. \) \( e_r.\text{in} = \bot \iff \exists \sigma'. \sigma' r_i = e_r \) \( \text{(F.3b)} \)

\( \forall \sigma, \sigma' \in \Sigma. \) \( \sigma \neq \sigma' \implies \sigma \sqsubseteq \sigma' \lor \sigma' \sqsubseteq \sigma \) \( \text{(F.4a)} \)

\( \forall \sigma \in \Sigma. \) \( \sigma r_i \sqsubseteq \sigma a_i \sqsubseteq \sigma b_i \) \( \text{(F.4b)} \)

\( \forall w_i, w_i', \sigma. \) \( w_i \xrightarrow{\text{fwd}} \sigma \land w_i' \xrightarrow{\text{fwd}} \sigma \implies w_i = w_i' \) \( \text{(F.5a)} \)

\( \forall \sigma. \) \( \mathcal{T}(\sigma b_i) \land \neg(\sigma r_i \xrightarrow{\text{fwd}} \sigma b_i) \implies \exists w_i, w_i' \xrightarrow{\text{fwd}} \sigma \) \( \text{(F.5b)} \)

\( \forall w_i, \sigma. \) \( w_i \xrightarrow{\text{fwd}} \sigma \implies w_i.a < \sigma b_i \land \neg(\sigma r_i \xrightarrow{\text{fwd}} \sigma b_i) \) \( \text{(F.5c)} \)

\( \forall w_i, \sigma. \) \( \sigma r_i \xrightarrow{\text{fwd}} \sigma a_i \land w_i \sqsubseteq \sigma r_i \implies w_i.a < \sigma a_i \) \( \text{(F.6)} \)

\( \forall w_i, w_i', \sigma. \) \( w_i \xrightarrow{\text{fwd}} \sigma \land w_i' a < \sigma r_i \implies w_i \neq w_i' \) \( \text{(F.7a)} \)

\( \forall w_i, w_i', \sigma. \) \( w_i \xrightarrow{\text{fwd}} \sigma \land w_i' \xrightarrow{r_i} \sigma \implies w_i \neq w_i' \) \( \text{(F.7b)} \)

Fig. 7. Signature representation of the properties of a snapshot data-structure that uses forwarding.

B in the multi-scanner case, where different virtual scans have different forwarding arrays. In the case of single-scanner algorithms, the virtual and abs scans coincide, and the forwarding array of each virtual scan is instantiated with the global forwarding array.

### 2.5.2 Forwarding Signature Formally

Our axiomatization of forwarding principles is given in Fig. 7. The signature, which we refer to as Forwarding, declares as inputs the sets of write events \( W_i \) (for each \( i \)), and the set of scan events \( S \). It then postulates the existence of a set of virtual scans \( \Sigma \), with a mapping \([-\]^* \) from abs to virtual scans, and a relation \( \text{fwd}_{\rightarrow} \) for visibility by forwarding that captures abstractly when a virtual scan observes a writer by forwarding. As before, these are abstract concepts, known to satisfy only the properties listed in the scope of \( \Sigma \). We will shortly provide the intuition for the axioms in Fig. 7, but first, let us enumerate how the axiomatization abstracts from Algorithm 1, so that the same signature applies to all three algorithms.

- The axiomatization does not assume that we have rep write events for forwarding \( (w_i b) \).
  Instead, it makes forwarding abstract by tying it to the relation \( \text{fwd}_{\rightarrow} \), allowing for multiple different forwarding methods.
- The axiomatization links virtual scans to abs scans by the function \([-\]^* \), to support multiple scanners. Physical scans can overlap in time, as long as the virtual scans do not.
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- The axiomatization does not assume any rep events operating with memory cell $X$, or for that matter, that such a cell even exists. Since $X$ is solely used to communicate to writers if and when to forward, we instead encode the conditions for forwarding in the axioms. This is necessary because the three algorithms decide differently on whether to forward or not.

2.5.3 Intuition Behind the Forwarding Signature. Out of the forwarding visibility relation $\text{fwd}_\rightarrow$, we define two additional visibility relations over abstract and virtual events: $\text{rf}_\rightarrow$ and $\text{wr}_\rightarrow$ as shown in Fig. 7. These relations define what we will use as an abstract notion of a virtual scan observing a writer directly or by forwarding ($\text{rf}_\rightarrow$), and a writer observing a prior write ($\text{wr}_\rightarrow$), respectively. In each algorithm, the $\text{fwd}_\rightarrow$ relation will be defined differently, but the definitions of $\text{rf}_\rightarrow$ and $\text{wr}_\rightarrow$ are constant in terms of $\text{fwd}_\rightarrow$. We next explain each of them.

$$
\text{rf}_\rightarrow w_i \sigma \triangleq (\text{wr}_a \text{rf}_a \sigma \land \sigma_r \text{rf}_r \sigma_b) \lor w_i \text{fwd}_\rightarrow \sigma
$$

Read visibility ($\text{rf}_\rightarrow$) captures that a virtual scan $\sigma$ either reads the value of $w_i$ by forwarding (disjunct $w_i \text{fwd}_\rightarrow \sigma$) or directly. The direct read requires that the scanner finds the value in $A[i]$ (conjunct $\text{wr}_a \text{rf}_a \sigma \land \sigma_r \text{rf}_r \sigma_b$) and that the scanner finds $\bot$ for the forwarded value in $B[i]$ (conjunct $\sigma_r \text{rf}_r \sigma_b$).

$$
\text{wr}_\rightarrow w_i \sim w_i' \triangleq w_i a \triangleq w_i' a
$$

Write visibility ($\text{wr}_\rightarrow$) orders the abs writers of a common memory cell, and is defined by the ordering of underlying rep writers $w_i a$. Ordering abs writers this way works because the rep write is the unique point where the abs writer communicates its value to the data structure. This holds for all three algorithms including the multi-writer variants.

We next explain the forwarding properties. These are divided into three groups, as shown in Fig. 7. The first group consists of the structural properties related to writers and scanners.

$$
\forall s \in \text{dom}([-]^*) . \quad s^* \subseteq s \quad \quad \text{(F.1 revisited)}
$$

$$
\forall i, s \in \mathcal{T}(S) . \quad \exists w_i . \quad w_i \text{rf}_a s^* \land w_i.\text{in} = s.\text{out}[i] \quad \quad \text{(F.2 revisited)}
$$

Property (F.1) states that each virtual scan must be inside the interval of the abs scan it represents. Property (F.2) states that a terminated abs scan $s$ must have a virtual scan representative that, for each index $i$, observed some write $w_i$ such that the value written by $w_i$ is the one the scan $s$ returned for $i$. Property (F.2) is similar to Property (S.1), and essentially says that the observing scan correctly reads the array. On the other hand, Property (F.1) is similar to a property of linearization points, whereby a linearization point must reside within the interval of the considered event. Here instead, we have a whole virtual scan $s^*$ representing when $s$ logically occurred.

$$
\forall e_r \in A[i].W . \quad \exists w_i . \quad w_i a = e_r \quad \quad \text{(F.3a revisited)}
$$

Property (F.3a) ensures that the only rep event that can write into $A[i]$ is $w_i a$ of some abs write $w_i$. This means that only abs writes into cell $i$ can have a rep event for writing into $A[i]$; scans and writes into a cell different from $i$ cannot, since if they did, they would have to be equal to some $w_i$ as per Property (ES.2) of event signatures. Also, an abs write can write into $A[i]$ at most once, since any other write must equal $w_i a$.

$$
\forall \sigma, e_r \in \sigma.B[i].W . \quad e_r.\text{in} = \bot \iff \exists \sigma'. \quad \sigma' r_i = e_r \quad \quad \text{(F.3b revisited)}
$$

The easiest way to explain Property (F.3b) is to consider its instantiation for Algorithm 1, where virtual scanners are abs scanners, which, moreover, all share the same array $B$.

$$
\forall e_r \in B[i].W . \quad e_r.\text{in} = \bot \iff \exists \sigma'. \quad \sigma' r_i = e_r \quad \quad \text{(F.3b')}
$$

The simplified Property (F.3b’) says that the only rep event that can write $\bot$ into $B[i]$ is $\sigma' r_i$ of some abs scan $\sigma'$. Thus, no event except scanners can write $\bot$ into $B[i]$, and the writing can be done at
most once. Additionally, the rep event $\sigma_i$ can only write $\bot$ into $B[i]$. Property (F.3b) generalizes Property (F.3b') by allowing for virtual scanners, each of which has their own $B$ array.

$$\forall \sigma, \sigma' \in \Sigma. \quad \sigma \neq \sigma' \implies \sigma \subset \sigma' \lor \sigma' \subset \sigma$$  
(F.4a revisited)

Property (F.4a) captures that virtual scans never overlap, and are thus totally ordered. This property holds for single-scanner algorithms by assumption, but must be proved in the multi-scanner case.

$$\forall \sigma \in \Sigma. \quad \sigma_i \subset \sigma a_i \subset \sigma b_i$$  
(F.4b revisited)

Property (F.4b) simply reflects that the rep events of initializing the scan for element $i$, reading $A[i]$, and then closing the scanning for element $i$ by reading $B[i]$, are invoked sequentially in the code of scan. This is readily visible in Algorithm 1 and remains true in Algorithms 2 and 3.

The second group are the structural properties of the forwarding visibility.

$$\forall w_i, w'_i, \sigma. \quad w_i \xrightarrow{\text{fwd}} \sigma \land w'_i \xrightarrow{\text{fwd}} \sigma \implies w_i = w'_i$$  
(F.5a revisited)

$$\forall \sigma. \quad \tau(\sigma b_i) \land \neg(\sigma_i \xrightarrow{\text{rf}} \sigma b_i) \implies \exists w_i. \ w_i \xrightarrow{\text{fwd}} \sigma$$  
(F.5b revisited)

$$\forall w_i, \sigma. \quad w_j \xrightarrow{\text{fwd}} \sigma \implies w_j a < \sigma b_i \land \neg(\sigma_i \xrightarrow{\text{rf}} \sigma b_i)$$  
(F.5c revisited)

Property (F.5a) captures that a scan can observe at most one forwarded write for a given index $i$. Property (F.5b) says that a forwarding of some write of cell $i$ will reach the virtual scan $\sigma$ if $\sigma b_i$ did not observe the writing of $\bot$ in $B[i]$ by $\sigma_i$. Contrapositively, $\sigma_i \xrightarrow{\text{rf}} \sigma b_i$ holds if no write of cell $i$ was forwarded to $\sigma$. Property (F.5c) states the dual implication direction of Property (F.5b) and that a forwarded write $w_i$ must have written into $A[i]$ by $w_i a$ before the scanner $\sigma$ read the forwarded value by $\sigma b_i$. This property captures that writers must first write their value directly before attempting to forward.

The third group are the properties that capture the forwarding principles of Jayanti. We have originally discovered these properties by extracting the common patterns from our proofs of the snapshot signature for the three algorithms, and only afterwards discovered that they actually correspond quite closely to Jayanti’s forwarding principles. To describe how our axioms capture the forwarding principles, we state the principles below in English, verbatim as Jayanti does, but using our notation. We also use virtual scans instead of abstract scans to make the connection to multi-scanner algorithm direct; Jayanti only stated the principles in terms of Algorithm 1. As we shall see, our axiomatization modifies the principles slightly to encompass Algorithms 2 and 3. We also highlight and number substatements so that we can relate them to our axioms in discussion.

1. Suppose that (i) a scan operation $\sigma$ misses a write operation $w_i$ writing $v$ because $\sigma$ reads $A[i]$ before $w_i$ writes in $A[i]$. If (ii) $w_i$ completes before $\sigma$ performs $s_{off}$, then (iii) $w_i$ will have surely informed $\sigma$ of $v$ by writing $v$ in $B[i]$.

2. Suppose that (iv) a scan operation $\sigma$ reads at $\sigma b_i$ a non-$\bot$ value in $B[i]$ by some write operation $w_j$. Then (a) (v) $w_j$ is concurrent with $\sigma$, and (b) If (vi) $w_j'$ is any write operation that is executed after $w_j$, then (vii) $w_j'$ completes only after $\sigma$ performs $s_{off}$.

We start with Property (F.6), which relates to Principle (1) as follows.

$$\forall w_i, \sigma. \quad \sigma_i \xrightarrow{\text{rf}} \sigma b_i \land w_i \sqsubseteq \sigma \implies w_i a < \sigma a_i$$  
(F.6 revisited)

Here, $\sigma_i \xrightarrow{\text{rf}} \sigma b_i$ corresponds to the negation of Statement (iii), $w_i \sqsubseteq \sigma$ corresponds to Statement (ii), and $w_i a < \sigma a_i$ corresponds to the negation of Statement (i), altogether combining into an equivalent, by contraposition, of Principle (1). We now explain each correspondence individually.

- $\sigma_i \xrightarrow{\text{rf}} \sigma b_i$ states that nothing was forwarded to $\sigma$, because the event $\sigma b_i$ of reading $B[i]$ observes the event $\sigma_i$ of clearing $B[i]$. This is a slightly stronger statement than the negation of Statement (iii), namely, that the write $w_i$ was not forwarded to $\sigma$. The difference between
these statements will be covered by Property (F.7b), which will handle the scenario where another write, not $w_i$, is forwarded to $\sigma$.

- $w_i \ll prec_{\sigma} \ll w'$, where $\ll prec_{\sigma}$ is the relational composition of $\ll$ and $\ll prec$, roughly translates to “$w_i$ terminates before some other writer (possibly writing to a different memory cell) that in turn was observed by $\sigma$”. We know such other writer exists, since the domain of $\ll prec$, the last relation in the composition, consists only of writers. While this looks nothing like Statement (ii), both statements imply that $w_i$ must have been forward if the writing to $A[i]$ was missed. We prove this with different proofs for each algorithm, with Lemma C.2 for Algorithm 1 and Lemma D.2 for Algorithms 2 and 3. We use the statement $w_i \ll prec_{\sigma} \ll w'$, $\sigma$, instead of potentially another statement involving $s_{off}$ (as Jayanti’s original statement does), because it makes the forwarding signature export fewer rep events, thus making it a bit more parsimonious.

- $w_i \ll prec_{\sigma} \ll w'$ states that $w_i a$ wrote to $A[i]$ before $\sigma a_l$ read from it, which directly corresponds to the negation of Statement (i).

For Principle (2a), we need to further deviate from Jayanti’s original formulation, since the latter is too specific to Algorithm 1, and does not scale as-is to the multi-writer algorithms. The point of this principle is to ensure that a writer may only forward values that are, intuitively, “current”. The principle itself ensures this property, but only in the single-writer case of Algorithm 1. Indeed, in Algorithm 1, a writer only forwards its own value (and other writes to the same pointer are prohibited by assumption); therefore, if a write is concurrent to the scan, then it can only forward a value that is current. But in a multi-writer case, as we shall see in Section 4, we want to allow a writer to forward values of other writers. In particular, a writer that writes and then stalls indefinitely could still be forwarded by another write to some concurrent scan. But we still need to ensure that the forwarded write is not arbitrarily old. We do so by insisting that a write $w_i$ forward to $\sigma$ cannot occur before a write $w_i'$ which wrote to $A[i]$ with $w_i a$ before $\sigma a_l$ started, i.e., executed $\sigma r_i$. In other words, we require that the forwarded write $w_i$ does not “happen before” any write $w_i'$ that started before the start of the scan. A write $w_i'$ has a newer value, so its existence should prevent $w_i$ from being forwarded. We formalize this as Property (F.7a).

$$\forall w_i, w'_i, \sigma. \quad w_i \xrightarrow{fwd} \sigma \land w'_i \ll prec_{\sigma} a_l \Rightarrow w_i \not\preceq w'_i \quad \text{(F.7a revisited)}$$

This captures that writers must forward to scan $\sigma$ either the latest write that wrote to $A[i]$ before $\sigma r_i$, or some write that occurred after $\sigma r_i$, thus ensuring that the forward write is current.

Lastly, Property (F.7b) relates to the contrapositive of Principle (2b).

$$\forall w_i, w'_i, \sigma. \quad w_i \xrightarrow{fwd} \sigma \land w'_i \ll prec_{\sigma} \ll w_i \Rightarrow w_i \not\preceq w'_i \quad \text{(F.7b revisited)}$$

Again, $w_i \xrightarrow{fwd} \sigma$ directly corresponds to Statement (iv). Formal statement $w'_i \ll prec_{\sigma} \ll w_i$ corresponds to the negation of Statement (vii), for the same reason as in Property (F.6), since (ii) is the negation of (vii). Finally, $w_i \not\preceq w'_i$ directly corresponds to the negation of (vi), with the exception that we use $\ll prec$ instead of Jayanti’s “executed after” to relate the writes $w_i$ and $w_i'$. The two correspond in the single-writer case, but our statement extends to the multi-writer case as well.

Algorithm 1 satisfies this set of properties of the Forwarding signature, and they are sufficient to prove that an algorithm is a snapshot algorithm according to the Snapshot signature, which is in turn sufficient to prove linearizability. In Section 3 we sketch why these claims holds, and in Section 4 we sketch how Algorithms 2 and 3 satisfy the Forwarding signature. The complete proofs are available in the appendix.
3 PROOF SKETCHES FOR ALGORITHM 1

We now sketch the proofs of our signatures implying linearizability, to illustrate the gist of reasoning by visibility. We begin by showing how histories that satisfy the Snapshot signature are linearizable. This proof is based on a similar proof by Chakraborty et al. [2015] for queues.

**Lemma 2.3.** Histories satisfying Snapshot signature (Fig. 5) are linearizable.

**Proof sketch.** Following Definition 2.2 of linearizability, we start by choosing a visibility relation $\prec$ that constructs the set $E_c = \{ e \mid \exists e' \in \mathcal{T}(E). e \prec e' \}$, and then extend it to a linearization $\prec$. Let $\prec$ be the visibility relation obtained by restricting the visibility relation postulated by Snapshot to $\mathcal{T}(S) \cup \bigcup_i \mathcal{W}_i$. Events outside of the restriction are non-terminated scans or writes that have not executed their effect yet; they do not modify the abstract state, and are hence not necessary for linearization. Additionally, by Property (S.5) we know that every terminated write is effectful. Thus going forward, we consider $E_c = \mathcal{T}(S) \cup \bigcup_i \mathcal{W}_i$.

We next extend the happens-before order $\prec$ on the selected set of events as follows. The idea is to keep extending this order with more relations between the events of the selected set until we reach a total order, which will be the desired linearization order. We first induce a helper order

$$\prec_w \triangleq (\prec \cup \text{wrDiff})^*$$

which ranges over writes (into different cells), where $w_i \xrightarrow{\text{wrDiff}} w'_j$ holds if $i \neq j$ and there exists a write $w_j$ and a scan $s$ such that $w_i, w_j \xrightarrow{\text{wr}} s$ and $w_j \prec w'_j$. Or, in English, we order $w_i$ before $w'_j$, if $w_i$ is observed by some scan along with $w_j$, and $w_j$ is ordered before $w'_j$. The $\prec_w$ order is a (strict) partial order by Properties (S.1) to (S.4) and (S.6) (see Lemma A.2 for the complete proof).

Next, we select an arbitrary total order $\prec_w$ over writes that extends $\prec_w$. Since the set of selected events is finite, such a total order always exists by Zorn’s Lemma. We use the order $\prec_w$ to determine which event we will consider as being last in the eventually constructed linearization order. We say that a write is latest-in-time if it is the greatest in $\prec_w$, while a scan is latest-in-time if it is maximal in $\prec$ and all the writes that it observes are the greatest in $\prec_w$ for their respective memory cell. As long as the set of events is non-empty, there will exist a latest-in-time event under this definition by Properties (S.1) to (S.4) (see Lemma A.4 for the complete proof). By inductively selecting the latest-in-time event, we construct a total order of events that is consistent with $\prec$ and snapshot semantics, and is thus a linearization order. More concretely, we add the latest-in-time event to the end of the linearization order, forming a chain with previously added events, and repeat this step with the set of events excluding the last selected.

Next, we show that histories satisfying the Forwarding signature also satisfy the Snapshot signature. Because some Forwarding properties use rep events, to be able to efficiently use them, we first need a lemma that turns a $\prec$ relation between virtual or abs events into a $\prec$ relation between rep events. The lemma will also be useful in showing that Algorithm 1 (and Algorithms 2 and 3 in Section 4) satisfy the Forwarding signature, once Properties (F.4b) and (F.5c) have been proven.

**Lemma 3.1 (Happens-before of subevents).** Assume Properties (F.4b) and (F.5c). For events $e, e' \in \bigcup_i W_i \cup \Sigma$ of the Forwarding signature (Fig. 7) where $e'$ is populated with at least one rep event from its event structure (Fig. 6), if $e \prec e'$ holds (where $\prec$ follows the definition from Fig. 7) then there exists $e_r$ and $e'_r$ where $e_r \prec e'_r$ holds, with $e_r$ and $e'_r$ belonging to $e$ and $e'$ respectively.

**Proof.** We can view $e \prec e'$ as a chain of $\prec_1$ relations connecting $e$ and $e'$; that is, $e \prec_1 \cdots \prec_1 e'$, where by definition $\prec_1 = (\subset \cup \prec) = (\subset \cup \xrightarrow{\text{wr}} \cup \xrightarrow{\text{rf}})$. It is easy to see that each abs event in this chain is populated with a rep event. Indeed, if an abs event’s relation $\prec_1$ to its successor is realized by $\subset$, then the event is terminated and all of its rep events are executed (and by our axioms, each abs
event must have at least one rep event). Alternatively, if the abs event’s relation $\prec_1$ to its successor is realized by $\prec$, then the event must be populated as per the definitions of reads-from and writing visibility in Fig. 7, and in the $\text{fwd}_r$ case by Property (F.5c). This leaves out $e'$, which has no successor, but $e'$ is populated by assumption.

Now the proof is by induction on the length of the above chain. The base case is when $e \prec_1 e'$. If $e \sqsubset e'$, the proof is by Property (RB.2) giving $e_i \sqsubset e'_j$ for arbitrary $e_i$ and $e'_j$ belonging to $e$ and $e'$ respectively. If $e \nless e'$, the proof is by the definition of $\nless$ in Fig. 7, giving us $w_ia \nless w'_ja$ with $e = w_j$ and $e' = w'_j$. For the last case, $e \nrightarrow e'$, we have $e = w_j$ and $e' = e'$ with either $w_ia \nrightarrow w'_ja \sigma' a_i$ or $w_j \nrightarrow w'_j \sigma'$. The former case directly exhibits rep events that satisfy our goal. The latter does so as well, since by Property (F.5c), it must be $w_ia \nless \sigma' b_j$. The inductive step is similar, with the addition that we need transitivity of $\prec$ and Property (F.4b) to join events.

**Lemma 3.2.** Histories satisfying Forwarding signature (Fig. 7) also satisfy Snapshot signature.

**Proof Sketch.** We start by determining the instantiations of $\text{wr}_j$, $\text{rf}_j$, and $\prec$ for the Snapshot signature. Since the effect of $w_j$ is observable only after the execution of $w_ia$, we let $w_j \in \text{wr}_j$ iff $w_ia$ is defined. For the other notions, we first define a helper visibility over scans $\text{sc}_j$, which we use to define $\prec$. The highlighted relations $\text{rf}_j$ and $\text{wr}_j$, originate from the Forwarding signature.

$w_j \text{rf}_j s \equiv w_j \text{rf}_j s^* \quad s \text{sc}_j s' \equiv s^* \sqsubset s'^*$ \quad \prec \equiv \text{rf}_j \cup \text{wr}_j \cup \text{sc}_j$

Property (V.1) holds since if $e \prec e'$ and $e' \sqsubset e$, then we can construct a $\equiv$-cycle of rep events with Lemma 3.1, contradicting irreflexivity of $\prec$ (Lemma 2.1). Property (S.1) holds directly by Property (F.2). Property (S.3) holds because by the definition of $w_j \text{rf}_j$, $\sigma$ in Forwarding signature, it is either $w_ia \text{rf}_j \sigma_j$, for which Property (M.3) ensures uniqueness, or $w_j \text{fwd}_i \sigma$ where Property (F.5a) ensures uniqueness. Property (S.4) of total ordering on effectful writes follows because in Forwarding signature, $w_j \text{wr}_j$ is defined as $w_ia \prec w'_ja$, and the rep writes are totally ordered by Property (M.4). Property (S.5) that all terminated writes are effectful holds because for every terminated write $w_j$, we must have $w_ia$ defined.

We next prove Property (S.2). We sketch the proof by assuming there is a write $w'_j$ such that $w_j \prec w'_j \prec s$ where $w_j \text{rf}_j s$, and then deriving a contradiction. Unfolding the definition of $\text{rf}_j$ in $w_j \text{rf}_j s$, we get $w_j \text{rf}_j \sigma$ where $\sigma = s^*$. Also, from $w'_j \prec s$, using Property (F.1), it turns out that $w'_j \prec \sigma$, and moreover, the latter splits into two possible cases: $w'_j \sqsubset \text{sc}_j \sigma$ and $w'_j \prec \sigma_j$. The full proof of how the cases follow from the premise and that they exhaust the possibilities is given in Appendix B. Note that the proposition $w'_j \sqsubset \text{sc}_j \sigma$ is the key part of Properties (F.6) and (F.7b), while $w'_j \prec \sigma_j$ is the key part of Property (F.7a); in fact, this exhaustion of cases of $w'_j \prec \sigma$ is what let us rediscover and mathematically formulate the forwarding principles. The proof now proceeds by analyzing $w_j \text{rf}_j \sigma$. From the definition of reads-from visibility in Fig. 7, we also get two cases: $(w_ia \text{rf}_j \sigma_j \wedge \sigma_j \text{sc} \sigma_j)$ and $w_j \text{fwd}_j \sigma$. For a total of four cases when combined with the above. In each case, contradiction follows easily, relying in different cases on a different subset of Properties (F.6), (F.7a), (F.7b) and (M.2).

Finally, we prove Property (S.6); that is, $w_i, w_j \text{rf}_j s$ and $w'_j, w'_j \text{sc}_j s'$ with $w_i \prec w'_j$ and $w'_j \prec w_j$ lead to contradiction. As before, we first compare the assumptions $w_i, w_j \text{rf}_j s$ and $w'_j, w'_j \text{sc}_j s'$ into $w_i, w_j \text{rf}_j \sigma$ and $w'_j, w'_j \text{sc}_j \sigma'$, where $\sigma = s^*$ and $\sigma' = s'^*$. Next, by property Property (F.4a), the virtual scans are totally ordered by $\sqsubset$, thus either $\sigma \sqsubset \sigma'$ or $\sigma' \sqsubset \sigma$ (the case when $\sigma = \sigma'$ contradicts reads-from uniqueness Property (S.3)). In the first case (the second is symmetric), we have $w'_j \prec w_j \text{rf}_j \sigma \sqsubset \sigma'$, and thus also $w'_j \prec w_j \prec \sigma'$. In other words, we have a scan $\sigma'$ observing a write $w'_j$ with another intervening write $w_j$ showing up in between. But this contradicts the argument we carried out for Property (S.2) above. □
LEMMA 3.3. Every execution of Algorithm 1 satisfies the Forwarding signature (Fig. 7).

Proof sketch. We instantiate the set of virtual scans to be the same as the set of abs scans (since Algorithm 1 is single-scanner), and define \([-\cdot]\) by \(s^* = s\). Property (F.1), requiring \(s^* \in s\), follows immediately because every event \(e\) satisfies \(e \subseteq e\). Since Algorithm 1 is a single-scanner algorithm, the total order of virtual scans Property (F.4a) naturally follows.

We instantiate forwarding visibility \(w_i \xrightarrow{\text{fwd}} s \doteq w_i \xrightarrow{\text{rf}} s_i\) to directly capture how forwarding is defined for Algorithm 1. Properties (F.2), (F.3a), (F.3b) and (F.4b) hold by the structure of the algorithm, while Properties (F.5a) to (F.5c) follows from the \(\text{fwd}\) instantiation (proved in Appendix C). For the remaining Properties (F.6) to (F.7b), we need two helper lemmas (also proved in Appendix C):

- Lemma C.1: \(s_{on} \xrightarrow{\text{rf}} w_i x\) iff we have \(s_{on} \xleftarrow{\text{rf}} w_i x\) and \(s_{off} \xrightarrow{\text{rf}} w_i x\).
- Lemma C.2: If \(w_i' \xrightarrow{\text{rf}} s\) then \(s_{off} \xrightarrow{\text{rf}} w_i' x\) and if \(w_i' \xrightarrow{\text{rf}} s\) was executed, then \(w_i' b \xleftarrow{\text{rf}} s_i\).

We focus here on proving Property (F.6) which involves showing if \(w_i' \xrightarrow{\text{rf}} s\) and \(s_i \xrightarrow{\text{rf}} s_i\) then \(w_i' a \xleftarrow{\text{rf}} s_i\). By Lemma C.2, we have \(s_{off} \xrightarrow{\text{rf}} w_i' b\) and \(w_i' b \xleftarrow{\text{rf}} s_i\) if \(w_i' b\) was executed. Consider if we have \(s_{on} \xleftarrow{\text{rf}} w_i' x\), then by Lemma C.1 we have \(s_{on} \xrightarrow{\text{rf}} w_i' x\) which in turn implies that \(w_i' b\) must have executed, however that contradicts Property (M.2) by \(w_i' b\) occurring in between \(s_i \xrightarrow{\text{rf}} s_i\), thus we must have \(s_{on} \xleftarrow{\text{rf}} w_i' x\). By Property (RB.1) over \(w_i' a \xleftarrow{\text{rf}} w_i' x\) and \(s_{on} \xleftarrow{\text{rf}} s_i\), since \(s_{on} \xrightarrow{\text{rf}} w_i' x\) contradicts \(s_{on} \xleftarrow{\text{rf}} w_i' x\), we must have \(w_i' a \xleftarrow{\text{rf}} s_i\), which is our goal.

4 MULTI-WRITER ALGORITHMS

To allow for correct non-blocking algorithms with multiple concurrent writers, as well as for multiple concurrent scanners in the case of Algorithm 3, Jayanti uses LL (load-link), SC (store-conditional) and VL (validate) operations [Jensen et al. 1987] in addition to simple memory mutation and dereference. We thus need to extend the register signature to account for the new operations. Additionally, Algorithms 2 and 3, both being multi-writer algorithms share more structure than what we capture in the Forwarding signature of Fig. 7. We thus introduce a new MWForwarding signature to capture the commonality of the multi-writer algorithms. We then show that histories satisfying MWForwarding also satisfy Forwarding and that Algorithms 2 and 3 satisfy MWForwarding.

4.1 LL/SC Registers

For a memory cell \(X\), LL(\(X\)), or load-link, reads from, and returns the value of \(X\). It also records the time of the read, for use in future SC and VL operations of the same thread. SC(\(X, v\)), or store-conditional, writes \(v\) into \(X\) and returns \(\text{true}\) if no other write into \(X\) occurred since the most recent LL(\(X\)) from the same thread. Otherwise, SC returns \(\text{false}\), keeping \(X\) unchanged. Finally, VL(\(X\)), or validate, returns truth values identically to SC, but does not mutate \(X\).

To capture the properties of memory with the LL, SC and VL operations, we introduce a new signature in Fig. 8, which we call LL/SC registers, or LLReg for short. It extends AReg from Fig. 3. We will treat each execution of LL, SC and VL as a read-like event, i.e., they will observe some write by \(\xrightarrow{\text{rf}}\). Additionally, we denote the LL visibility \(\xrightarrow{\text{ll}}\) to link LL events with their corresponding SC/VL events. To capture the concept of SC/VL operations being successful, we introduce the set of successful events \(S\). If an SC event is in \(S\), it was successful and must have written to memory, therefore we treat it as a write-like event. These sets are formally defined in the top part of Fig. 8.

Further, referring to Fig. 8, the Properties (M.1a) to (M.4) are almost identical to (M.1) to (M.4) from AReg from Fig. 3. The main difference is that we substitute quantification over writes and reads with quantification over write-like events \((W_c = W \cup (SC \cap S))\) and read-like events \((R_c = R \cup LL \cup SC \cup VL)\), respectively. We also split Property (M.1) into three: (M.1a) says that each terminated read-like event has some observed write-like event; (M.1b) says that the output of
reads and LL is the same as the input of the observed write-like event, and (M*1c) says that the output of SC and VL depends on its success. Additionally, Property (M*5) states that each SC and VL event has a terminated LL event it is related to in \( \ll \). Property (M*6) says that each \( \ll \) pair has the same abs parent, capturing that the two events of the pair were executed in the same thread (as per our assumption in Section 2.2.1 that each event is single-threaded), and that no other LL event occurs in between the pair. Finally, Property (M*7) captures that successful events observe the same write-like event as their \( \ll \)-related LL event, meaning no other modifications occurred between the LL event and the successful event.

There are several important properties of LL/SC/VL that the algorithms relies upon for synchronization. We codify them in the following three lemmas and illustrate in Fig. 9. For each of the lemmas, the events \( l, c, w \), and variants, operate over the same memory cell \( X : LLReg \).
Lemma 4.1. Let \( l \Downarrow\ c \) and \( l' \Downarrow\ c' \) with \( c, c' \in \mathcal{T}(SC \cap S) \). If \( c < c' \) then \( c < l' \).

Proof. By Property (M’.5), we know that \( l \) and \( l' \) are terminated, and by Property (M’.1a), we have that each of \( l, l', c \) and \( c' \) observes some write-like event. By Property (M’.7) with \( c, c' \in S \), then \( l \) and \( c \) observe the same write-like event, and dually for \( l' \) and \( c' \), i.e., there exists \( w \) and \( w' \) such that \( w \Downarrow\ l \) and \( w \Downarrow\ c \) and \( w' \Downarrow\ l' \) and \( w' \Downarrow\ c' \). By Property (M.4), we either have \( w' < c \) or \( c < w' \). In the former case, we have \( w' < c < c' \), which contradicts Property (M.2) by there being a \( c \) in between \( w' \Downarrow l' \). Thus, we must have the latter case, giving us \( c \Downarrow w' \), letting us construct \( c \Downarrow w' \Downarrow l' \) which implies our goal \( c < l' \).

Lemma 4.2. Let \( l \Downarrow\ c \) with either \( c \in \mathcal{T}(SC) \) or \( c \in \mathcal{T}(VL \setminus S) \). Then there exists some \( w \in W_c \) such that \( w \Downarrow l \) and \( w \Downarrow \).

Proof. If \( c \in \mathcal{T}(SC \cap S) \), then \( c \) must have successfully written; thus \( c = w \) trivially satisfies \( w \Downarrow c \) and \( w \Downarrow \). The latter holds, because otherwise we get \( c < l \) and \( l \Downarrow c \) which contradicts irreflexivity of \(<\). Next, consider \( c \notin S \), where \( c \in \mathcal{T}(SC \cup VL) \).

By Property (M’.5), \( l \) is terminated, and by Property (M’.1a), for \( l \) and \( c \) there exist write-like events \( w, w' \in W_c \) observed by \( l \) and \( c \), respectively. By Property (M’.7), these write-like events must be distinct, because otherwise \( c \) would have been successful. In other words, \( w \Downarrow\ l \) and \( w' \Downarrow\ c \) where \( w \neq w' \). Taking \( w' \) to be the required write of the lemma, we trivially have \( w' \Downarrow c \). To show that also \( w' \Downarrow l \), we assume \( w' \Downarrow l \) and derive a contradiction. By Property (M.4), we either have \( w < w' \) or \( w < w' \). In the first case, we contraduct Property (M.2) for \( w \Downarrow l \) by \( w < w' \Downarrow l \).

In the second case, we contradict Property (M.2) for \( w' \Downarrow l \) by \( c < w' \Downarrow l \).

Lemma 4.3. Let \( l \Downarrow\ c \subset l' \Downarrow\ c' \) where \( c, c' \in \mathcal{T}(SC) \). If within the time frame of the events \( l, l', c, \) and \( c' \) there are no mutations to register \( X \) except by \( SC \), then there exist some \( l'' \) and \( c'' \) (over \( X \)) with \( l'' \Downarrow\ c'' \) such that \( l'' \Downarrow l \) and \( c'' \Downarrow c' \).

Proof. From the assumption, by Lemma 4.2, there exists some \( c_0 \) and \( c_0' \) such that \( c_0 \Downarrow l \) and \( c_0 \Downarrow c \) and \( c_0 \Downarrow l' \) and \( c_0' \Downarrow c' \). Since there are no other writes, we can only have \( c_0, c_0' \in SC \cap S \), by Property (M’.5) there has to be some \( l_0' \) such that \( l_0' \Downarrow c_0' \). Let \( l_0' \) and \( c_0' \) be the existentials, we trivially have \( c_0 \Downarrow c' \) and \( c_0' \in SC \cap S \). For \( l_0' \Downarrow l \), assume we have \( l_0' \Downarrow l \), we derive a contradiction.

We must have \( c_0 < c_0' \), since we otherwise we contradict \( c_0 < l' \), then \( c_0 < l_0' \Downarrow l' \), meaning we also must have \( c_0 < l_0' \) for \( c_0' \) to be successful, however this contradicts \( c_0 \Downarrow l \) by \( c_0 < l_0' < l \).

Lemma 4.1 says that if we have two successful LL/SC pairs, then the intervals of the two pairs do not overlap. To see the intuition behind the other two lemmas, consider the special case when the order \( < \) is total (e.g., it is a linearization order), so that \( \Downarrow\ ) is equivalent to \( \sqsupset \). In that case, Lemma 4.2 states that an LL/SC interval, or an LL/VL interval with a failing VL, must contain some successful write-like event. Similarly, Lemma 4.3 states there exists a successful LL/SC interval contained within the interval defined by two sequential LL/SC intervals, even if the latter two intervals themselves correspond to failing SC’s. The last property will be used by multi-writer algorithms as follows. Referring to Algorithm 2, a write procedure will try to execute forwarding by invoking two sequential intervals of LL/SC pairs into the cell \( B[i] \). Even if both SC’s fail, Lemma 4.3 guarantees that another concurrent write will have invoked a successful LL/SC interval, thus forwarding the same value on behalf of the original write.

4.2 Multi-Writer Forwarding Signature

4.2.1 Description of Algorithm 2. We first discuss Algorithm 2, which will motivate the definition of MWForwarding signature. The common aspect with Algorithm 1 is that the writer communicates its value to the data structure by writing into \( A[i] \). Whereas Algorithm 1 may try to communicate the
which are new and correspond to a potential observer of

and

written by another overlapping write). This is the key property facilitating linearizability.

Visibility Reasoning for Concurrent Snapshot Algorithms

Algorithm 2. Jayanti’s multi-writer, single-scanner snapshot algorithm.

| resource A : array[n] of val |
| resource B : array[n] of val ∪ {⊥} |
| resource X : B := false |

1. **write(i : N, v : val) △

2. A[i] := v
3. x ← LL(X)
4. if x then
5. forw(i, x)
6. forward(i, x)
7. forward(i : N) △
8. LL(B[i])
9. v ← A[i]
10. x’ ← VL(X)
11. if x then SC(B[i], v)
12. scan: array[n] of val △
13. for i ∈ {0, ..., n - 1} do
14. B[i] := ⊥
15. X := true
16. for i ∈ {0, ..., n - 1} do
17. a ← A[i]
18. V[i] := a
19. X := false
20. for i ∈ {0, ..., n - 1} do
21. b ← B[i]
22. if b ≠ ⊥ then V[i] := b
23. return V

Algorithm 2. Jayanti’s multi-writer, single-scanner snapshot algorithm.

| event signature WRITE_i+ △ |
| resource A : array[n] of ARreg |
| resource X : LLReg |

| event signature VSCAN+ △ |
| B : array[n] of LLReg |
| r_i : B[i].W |
| a_i : A[i].R |
| b_i : B[i].R |
| on : X.W_c |
| off : X.W_c |
| f1, f2 ∈ FORWARD_i |

| event signature FORWARD_i △ |
| B_ceil : LLReg |
| b : B_ceil.LL |
| a : A[i].R |
| x : X.VL |
| b : B_ceil.SC |

Fig. 10. Resources and event signatures for abs writes, virtual scans and virtual forwarding corresponding to MWForwarding. The rep events correspond to the rep events with the same suffix in Algorithm 2, with the exception of on and off which are new and correspond to a potential observer of on and off respectively.

same value again by forwarding to B[i], in Algorithm 2, the forwarding procedure reads whichever value is currently present in A[i] and attempts to forward it to B[i] by means of LL and SC. Because multiple writes may be racing on A[i], the forwarding procedure may read and forward a different value from A[i] than the one the writer initially wrote. The forwarding procedure may even fail to forward anything. Nevertheless, Lemma 4.3 provides a guarantee that there will exist some successful forwarding among the concurrent processes, forwarding a value that is current (i.e., written by another overlapping write). This is the key property facilitating linearizability.

Other than that, the scan procedure is mostly the same in Algorithm 2 compared to Algorithm 1; the main difference is that the order between s_on and each sr_i is swapped, which is done in preparation for Algorithm 3. In Algorithm 1 this swap would result in a bug, causing interference between forwarding and initialization of B; the writes may fail to forward their value because the forwarding flag is turned on too late, or scans may erase forwarded values by the initializations.
signature MWForwarding $\triangleq$ Multi-Writer snapshot with forwarding

| Let               | Multi-Writer snapshot with forwarding |
|-------------------|--------------------------------------|
| $W_i$             | Set of all writes of cell $i$ of the snapshot in history with each $w_i \in \text{WRITE}_i$+ |
| $S$               | Set of all scans of the snapshot in history |
| $F_i$             | Set of all forwarding events in history with each $f \in \text{FORWARD}_i$ |
| $\Sigma$          | Set of virtual scans with each $\sigma \in \text{VSCAN}+$ |
| $[-\cdot]^*$      | Partial mapping of abs scans into virtual scans |
| $[-\cdot]_w$      | Write which executed the forwarding event |
| $[-\cdot]_f$      | Virtual scan observed by forwarding event |

\[
\forall s \in \text{dom}([-\cdot]^*). \quad s^* \subseteq s \quad \text{(F.1)}
\]

\[
\forall i, s \in \mathcal{T}(S). \quad \exists w_i, w_i^\prime, s^* \wedge w_i, \text{in} = s, \text{out}[i] \quad \text{(F.2)}
\]

\[
\forall e_r \in A[1].W. \quad \exists w_i. w_i.a = e_r \quad \text{(F.3a)}
\]

\[
\forall \sigma, e_r \in \sigma.A.B[i].W. \quad e_r.\text{in} = \bot \iff \exists \sigma'. \sigma' \ll = e_r \quad \text{(F.3b)}
\]

\[
\forall \sigma, e_r \in \sigma.A.B[i].W. \quad e_r.\text{in} \neq \bot \iff \exists f. f.b = e_r \quad \text{(F.3c)}
\]

\[
\forall \sigma, e_r \in \sigma.X.R. \quad \sigma.\text{on} \ll = e_r \iff \sigma.\text{on} \ll = e_r \wedge \sigma.\text{off} \ll = e_r \quad \text{(F.3d)}
\]

\[
\forall \sigma, \sigma' \in \Sigma. \quad \sigma \neq \sigma' \implies \sigma \ll = \sigma' \vee \sigma' \ll \sigma \quad \text{(F.4a)}
\]

\[
\forall \sigma \in \Sigma. \quad \sigma.r_i \ll = \sigma.\text{on} \ll \sigma.\text{on} \ll \sigma.\text{off} \ll \sigma.\text{off} \ll \sigma.b_i \quad \text{(F.4b)}
\]

\[
\forall w_i \in W_i. \quad w_i.a \ll = w_i.x \ll = w_i.f_i \ll = w_i.b \quad \text{(F.4c)}
\]

\[
\forall f \in F_i. \quad f.b \ll = f.a \ll = f.x \ll = f.b \ll = f.b \ll = f.b \quad \text{(F.4d)}
\]

\[
\forall f, w_i, \sigma. \quad w_i = [f]_w \wedge \sigma = [f]_\sigma \implies \sigma.\text{on} \ll = f.\text{on} \ll = f.\text{on} \ll = f.\text{off} \ll = \sigma.\text{off} \ll = f \quad \text{(F.5a)}
\]

\[
\forall f \in F_i. \quad \text{def}(f.b) \implies \exists \sigma. \sigma = [f]_\sigma \wedge \sigma.\text{on} \ll = f \quad \text{(F.5b)}
\]

Fig. 11. Signatures for multi-writer snapshot data-structure with LL/SC forwarding (Algorithms 2 and 3).

But this is safe to do in Algorithm 2 because the use of LL, SC and VL over $X$ and $B$ ensures that no forwarding (intended for a scan $s'$ prior to $s$) can succeed in-between $sr_i$ and $s_{on}$ of $s$. This is substantiated by the following lemma.

**Lemma 4.4.** Let $s'$ be a scan prior to $s$, i.e., $s' \ll = s$. If $f_x$ succeeds observing the same $s'_{on}$ that the write-like event $w$ preceding $f_x$ observed, i.e., $s'_{on} \ll = f_x$ and $s'_{on} \ll = w$, then $f_b$ must fail to write.

**Proof.** We assume that $f_b$ is successful, and derive contradiction. By Property (M.1b), there exists some write-like event $w$ such that $w \ll = f_b$, and by $f_b$ being successful, Property (M.7) implies that $f_b$ observes the same write-like event as $f_b$, i.e., $w \ll = f_b$. By Property (M.4), we either have $sr_i \ll = w$ or $w \ll = sr_i$, in the latter case we have $w \ll = sr_i \ll = f_b$, contradicting Property (M.2) by $sr_i$ being in between $w \ll = f_b$, thus we can only have $sr_i \ll = w$. Since Algorithm 2 is single-scanner, it follows that $s$ only starts after $s'$ has finished, thus we can derive the following:

\[
{s'_{on} \ll = s'_{off}} \ll = w \ll = f_b \ll = f_x
\]

which implies that we have $s'_{off}$ in between $s'_{on} \ll = f_x$, contradicting Property (M.2). \qed
4.2.2 The MWForwarding Signature. The rep-event signatures for multi-writer algorithms in Fig. 10 correspond closely to the rep events of Algorithm 2. Fig. 10 extends Fig. 6, adding the LL/SC register \( X \) and extending the event signatures. WRITE\(_{i}^{+} \) extends WRITE\(_{i} \) by events \( w_{a}, w_{f_{1}} \), and \( w_{f_{2}} \) from Algorithm 2, along with an event signature FORWARD, corresponding to the procedure FORWARD from Algorithm 2. VSCAN\(_{+} \) extends VSCAN by updating \( B \) to be an array of LL/SC registers, and adding four additional rep events: \( \sigma_{on} \) and \( \sigma_{off} \) corresponding to \( s_{on} \) and \( s_{off} \) from Algorithm 2, and \( \sigma_{on} \) and \( \sigma_{off} \), which are providing support for multi-scanner Algorithm 3. As we shall see in Section 4.3, in Algorithm 3, \( \sigma_{on} \) might not return before \( \sigma_{a} \), and similarly for \( \sigma_{off} \) and \( \sigma_{b} \). Thus, we need events that observe \( \sigma_{on} \) and \( \sigma_{off} \) returning before \( \sigma_{a} \) and \( \sigma_{b} \) respectively.

For Algorithm 2, we simply take \( \sigma_{on} = \sigma_{on} \) and \( \sigma_{off} = \sigma_{off} \).

Signature MWForwarding in Fig. 11 captures the common structure of Algorithms 2 and 3. Some aspects are preserved from Algorithm 1; for instance Properties (F.1), (F.2), (F.3a), (F.3b) and (F.4a) are the same as in Forwarding signature of Fig. 7. One difference is that we now have sets of forwarding events \( F \) that are instances of the forwarding procedure. The new signature introduces operations over such forwardings: \( [-]_{w} \) maps a forwarding to the write that invoked it, and \( [-]_{r} \) maps a forwarding to the virtual scan it is forwarding to. Forwarding visibility \( \text{fwd} \) is also updated: it now says that a forwarding reads a value of a write from \( w_{a} \) and RELAYs this value by \( f_{b} \).

Among the new properties, (F.3c) encodes that only \( f_{b} \) performs forwarding writes into \( \sigma_{B}[i] \). Property (F.3d) ensures that nothing wrote to \( X \) other than \( \sigma_{off} \) directly after \( \sigma_{on} \). Properties (F.4b) to (F.4d) corresponds to the rep event order of scans, writes and forwarding respectively. Lastly, Property (F.5a) ensures that every forwarding has a virtual scan to forward to, while Property (F.5b) states that \( f_{b} \) can only be performed if \( x_{a} \) observed \( \sigma_{on} \).

**Lemma 4.5.** Histories satisfying MWForwarding signature also satisfy Forwarding signature.

**Proof Sketch.** We instantiate \( \Sigma \) and \( [-]^{*} \) of Forwarding to be the same as the corresponding instantiations of MWForwarding. Properties (F.1), (F.2), (F.3a), (F.3b) and (F.4a) are shared between the signatures, therefore they trivially hold.

We next focus on proving Property (F.7b), leaving the remaining properties for Appendix D. Proving Property (F.7b) requires showing that \( w_{i} \overset{\text{fwd}}{\sigma} \) and \( w_{i} \overset{=}{{}\overset{\longrightarrow}{\text{rf}}} \sigma \) and \( w_{j} \overset{=}{{}\overset{\longrightarrow}{\text{rf}}} \sigma \) derive a contradiction. To show this, we need a helper lemma, proved in Appendix D (Lemma D.2): If \( w_{i} \overset{=}{{}\overset{\longrightarrow}{\text{rf}}} \sigma \) then surely \( w_{i} \) will be able to forward to \( \sigma \) in time, i.e., we have \( \sigma_{off} \not\overset{x}{\sigma} w_{a} \) and \( \sigma_{off} \not\overset{x}{\sigma} f_{a} \) and \( f_{b} \overset{}{\sigma_{a}} \sigma_{b} \) for every \( f \) executed by \( w_{i} \). From \( w_{i} \overset{\text{fwd}}{\sigma} \sigma_{i} \), by the definition of \( \text{fwd} \), we have \( w_{a} \overset{\text{rf}}{\sigma_{i}} f_{a} \) and \( f_{b} \overset{\text{rf}}{\sigma_{i}} \sigma_{b} \) for some \( f \). By Lemma 3.1, \( w_{i} \overset{=}{{}\overset{\longrightarrow}{\text{rf}}} \sigma \) implies \( w_{i} \overset{=}{{}\overset{\longrightarrow}{\text{rf}}} \sigma \). We can also infer (full proof in Appendix D) that \( \sigma_{on} \overset{}{\sigma} w_{a} \).

By Lemma D.2, we have \( \sigma_{off} \not\overset{x}{\sigma} w_{a} \), thus by Property (F.3d) we derive \( \sigma_{on} \overset{f_{a}}{\sigma} w_{a} \), implying \( w_{f_{1}} \) and \( w_{f_{2}} \) will be executed, let \( f' = w_{f_{1}} \) and \( f'' = w_{f_{2}} \). By Lemma D.2, we have \( \sigma_{off} \not\overset{x}{\sigma} f_{a} \) and \( \sigma_{off} \not\overset{x}{\sigma} f_{a} \) and \( f_{a} \overset{\sigma_{a}}{\sigma} f_{a} \), similarly above with Property (F.3d) we derive \( \sigma_{on} \overset{f_{a}}{\sigma} f_{a} \) and \( \sigma_{on} \overset{f_{a}}{\sigma} f_{a} \), meaning \( f_{b} \) and \( f_{b} \) will be executed, meaning we have two consecutive LL/SC pairs of \( f' \) and \( f'' \). Since we have \( \sigma_{on} \overset{f_{a}}{\sigma} w_{a} \) and \( f_{b} \overset{\sigma_{a}}{\sigma} \sigma_{b} \), it is not possible for any \( \sigma_{r} \) to occur in the intervals of the LL/SC pairs, therefore we can apply Lemma 4.3, thus there exists some \( f''' \) such that \( f''' \not\overset{x}{\sigma} f' \) and \( f''' \not\overset{x}{\sigma} f'' \).

By Property (M.4), we either have \( f_{b} \overset{\sigma_{b}}{\sigma} f''' \) or \( f''' \overset{\sigma_{b}}{\sigma} f_{b} \). In the first case, we have

\[
\begin{align*}
f_{b} \overset{\sigma_{b}}{\sigma} f''' \overset{\sigma_{b}}{\sigma} f_{b},
\end{align*}
\]

which contradicts Property (M.2) by \( f''' \) occurring in between \( f_{b} \overset{\sigma_{b}}{\sigma} \sigma_{b} \). In the second case, if we have \( f''' \overset{\sigma_{b}}{\sigma} f_{b} \) then by Lemma 4.1, and \( f_{b} \) and \( f''' \overset{\sigma_{b}}{\sigma} f_{b} \) being successful, we have \( f''' \overset{\sigma_{b}}{\sigma} f_{b} \), and this also follows in the case of \( f''' = f_{b} \). By Property (RB.1) over \( w_{a} \overset{f_{b}}{\sigma} f_{b} \) and \( f''' \overset{f''' \sigma_{b}}{\sigma} f''' \) we
resource \( A \): array\([n]\) of val
resource \( A_p \): array\([n]\) of val
resource \( B_p \): array\([n]\) of val \( \cup \{\perp\} \)
phase: \( \{1, 2, 3\} := 1 \)
resource \( X \): \( p_B p_A : \text{PID} \)
\( \text{sync} : B := \text{false} \)
\( \text{img} : \text{array}[n] \) of val
resource \( SS \): \( \text{sync} : \mathbb{B} := \text{false} \)

1: \( \text{write}(i : \mathbb{N}, v : \text{val}) \triangleq \)
2: \( A[i] := v \)
3: \( x \leftarrow \text{LL}(X) \)
4: if \( x.\text{phase} = 2 \)
5: \( \text{forward}(i, x.p_B) \)
6: \( \text{forward}(i, x.p_B) \)
7: \( \text{forward}(i : \mathbb{N}, p : \text{PID}) \triangleq \)
8: \( \text{LL}(B_p[i]) \)
9: \( v \leftarrow A[i] \)
10: \( x' \leftarrow \text{VL}(X) \)
11: if \( x' \) then \( \text{SC}(B_p[i], v) \)
12: \( \text{scan}: \text{array}[n] \) of val \( \triangleq \)
13: \( p \leftarrow \text{getPID} \)
14: \( \text{pushVS}(p) \)
15: \( \text{pushVS}(\text{PID}) \)
16: \( s_s \leftarrow SS \)
17: return \( s_s.\text{img} \)
18: \( \text{pushVS}(p : \text{PID}) \triangleq \)
19: \( x \leftarrow \text{LL}(X) \)
20: if \( x.\text{phase} = 1 \)
21: for \( i \in \{0 \ldots n - 1\} \) do
22: \( B_p[i] := \perp \)
23: \( \text{SC}(X, (2, x.p_A, p, x.\text{sync})) \)
24: \( x \leftarrow \text{LL}(X) \)
25: if \( x.\text{phase} = 2 \)
26: for \( i \in \{0 \ldots n - 1\} \) do
27: \( a \leftarrow A[i] \)
28: \( A_p[i] := a \)
29: \( \text{SC}(X, (3, p, x.p_B, x.\text{sync})) \)
30: \( x \leftarrow \text{LL}(X) \)
31: if \( x.\text{phase} = 3 \)
32: for \( i \in \{0 \ldots n - 1\} \) do
33: \( b \leftarrow B_{x.p_B}[i] \)
34: if \( b \neq \perp \)
35: \( V[i] := b \)
36: else
37: \( a \leftarrow A_{x.p_B}[i] \)
38: \( V[i] := a \)
39: \( ss \leftarrow \text{LL}(SS) \)
40: \( x' \leftarrow \text{VL}(X) \)
41: if \( ss.\text{sync} = x.\text{sync} \land x' \) then
42: \( \text{SC}(SS, \langle V, \neg ss.\text{sync} \rangle) \)
43: \( \text{SC}(X, (1, x.p_A, x.p_B, \neg x.\text{sync})) \)

Algorithm 3. Jayanti’s multi-writer, multi-scanner snapshot algorithm.

have \( w_i a \subset f”’b \subset f’b \) since we cannot have \( f”’b \subset f’b \) by \( f”’b \neq f’b \). Thus, we have

\( w_i a \subset w_i a \subset f”’b \subset f’b \subset f a \)

which contradicts Property (M.2) by \( w_i a \) occurring in between \( w_i a, f_i, f a \).

\( \square \)

**Lemma 4.6.** Every execution of Algorithm 2 satisfies the MWForwarding signature (Fig. 11).

**Proof.** Since this algorithm is single scanner, we can simply define the set of virtual scans to be the same as the set of abs scans and map each abs scan to itself, i.e., \( \Sigma = S \) and \( s^* = s \). Each rep event directly corresponds to their equivalent variant in Algorithm 2, except for \( \sigma_{on} \) and \( \sigma_{off} \), which are set as \( \sigma_{on} = \sigma_{on} = s_{on} \) and \( \sigma_{off} = \sigma_{off} = s_{off} \). Property (F1) holds by each event being a subevent of itself (\( e \subseteq e \)) and Property (F3a) holds since the algorithm is single-scanner. Properties (F2) to (F3c) and (F4b) to (F5b) holds directly by the structure of the algorithm and Property (F3d) can be proven by Lemma C.1, the same lemma that established this property for Algorithm 1, since the relative structure of \( s_{on} \) and \( s_{off} \) is the same.

\( \square \)

**4.3 Multi-Writer, Multi-Scanner**

Algorithm 3 is the final of the Jayanti’s snapshot algorithms, allowing unconstrained write and scan operations. It differs from Algorithm 2 in the implementation of scan, while write is mostly
unchanged. While Algorithm 3 allows for multiple concurrently running physical scanners, under the hood, at most one virtual scanner can be logically viewed as running at any point in time. The physical scans maintain their own local snapshot and forwarding arrays, but the shared resource X, via the fields X.p_A and X.p_B, keeps track which local snapshot array and local forwarding array, respectively, should be considered as the current data of the virtual scan. The physical scanners race to modify these fields in order to promote their own local arrays as the arrays of the virtual scan. However, X can be modified only by SC; thus, by Lemma 4.1, mutations to X occur sequentially, and by Lemma 4.2, at least one mutation exists. The sequential nature of these mutations justifies viewing them as belonging to a single, albeit virtual, scan. Once the virtual scan terminates, a snapshot is copied into SS (v_{ss}), which scanners read (s_{ss}) and return. We follow Jayanti in assuming that SS, even though physically a compound object, can still be considered an LL/SC register; this is not a loss of generality, as multi-word LL/SC registers exist [Jayanti and Petrovic 2005].

The pushVS procedure (short for “push virtual scan”) propels the virtual scanner to its completion, producing a snapshot in SS. The procedure is divided into three phases, tracked by the field X.phase, corresponding to the three phases of the previous algorithms: the first phase sets each memory cell in B to ⊥ with vr_i; the second phase reads the main memory A with va_i; and the third phase reads the forwarding memory B with vb_i. However, unlike before when B was a global array, here we have one B_p for each scan process p. A scanner that successfully writes with v_on in the first phase will have its B_p used for forwarding in the current virtual scan. For the second phase, the algorithm uses A_p to store the results of reading A, and similarly the process p which successfully writes with v_off will have its A_p used for the next phase. The third phase uses A_p and B_q from the processes stored in X to construct the snapshot, which is then stored in SS with v_{ss}.

Each phase starts with an LL(X) (three lines labeled vx) and ends with an SC(XX). This ensures that if multiple physical scanners run a phase simultaneously, only a single SC operation succeeds writing to X, committing the results of the phase. Thus, considered across all physical scanners, successful phases cannot overlap, by Lemma 4.1. Also, phases must run consecutively, because a phase in some scanner can only start if the previous phase, maybe executed by some other scanner, has been terminated, incrementing X.phase to enable the next phase (v_on, v_off, and v_end). Because the phases cannot overlap, must be executed in order, and correspond to the scanner phases in the previous algorithms, a trace of a virtual scan in Algorithm 3 has essentially the same structure as a trace of a physical scan in the Algorithms 1 and 2.

To ensure that SS is only mutated once per virtual scan, we have two toggle bits X.sync and SS.sync as part of X and SS respectively. These function such that every time v_end successfully writes, X.sync is negated. SS can only be written to by v_{ss}, which can only be executed if X.sync and SS.sync are equal1 and X did not change during the phase. Once v_{ss} successfully writes, SS.sync will be negated, meaning no more writes into SS can occur in this phase since X.sync and SS.sync are now distinct. We also know that there has to be exactly one successful write with v_{ss} before v_end, since v_{ss} only fails if some other v_{ss} succeeded, or some other v_end ended the phase.

It is also vital for linearizability that the virtual scan is in the time interval of the abs scan returning it. To ensure the virtual scan ends before the abs scanner ends, the latest virtual scan is finished before pushVS terminates. Dually, to ensure that the virtual scan was not started before the abs scanner starts, pushVS is executed twice; once to finish the currently running virtual scan, and once to generate a virtual scan that can be used by the abs scanner. Executing pushVS twice is similar to the idea of executing forward twice in Algorithm 2.

**Lemma 4.7.** Every execution of Algorithm 3 satisfies the MWForwarding signature (Fig. 11).

---

1Jayanti’s original presentation is dual, as the sync bits need to be distinct for writing to SS, but this is equivalent.
Proof sketch. We instantiate the write and forward rep events for the event structures of Fig. 10 with rep events of the same name. The structure of write and forward ensures that Properties (F.3a), (F.3c) and (F.4c) to (F.5b) are satisfied.

We instantiate virtual scans as a logical object corresponding to the rep events executed by pushVS constructing a complete virtual scan. The interval of a virtual scan $\sigma$ is instantiated as the smallest interval containing all its rep events. More formally, if we have $v, v'$ and $v''$ such that $v_{on} \xrightarrow{rf} v_{x}$ and $v'_{off} \xrightarrow{rf} v_{x}''$ and $v''_{ss}$ wrote to SS, then there exists a virtual scan $\sigma$. Its rep events are instantiated using the rep events of $v, v'$ and $v''$ as follows and instantiate $s^* = \sigma$ iff $v''_{ss} \xrightarrow{rf} s_{ss}$.

$\sigma_i \xrightarrow{\sigma_i} v_{on} \xrightarrow{\sigma_0} v_{x}' \xrightarrow{\sigma_a} v_{a}' \xrightarrow{v_{off}} v_{x}'' \xrightarrow{\sigma_{off}} v_{x}'' \xrightarrow{\sigma_{end}} v_{end} \xrightarrow{\sigma_{off}} v_{b}''\xrightarrow{\sigma_{off}} v_{b}''$

We also let $\sigma_{init} \xrightarrow{\sigma}_{ss} v_{x}'$ and $\sigma_{ss} \xrightarrow{\sigma_{ss}} v_{x}''$ and $\sigma_{end} \xrightarrow{\sigma_{end}} v_{end}'$, where $v_{end}'$ was successful and $v_{off} \xrightarrow{rf} v_{end}'$, however these do not need to be in the interval of $\sigma$. Properties (F.2), (F.3b), (F.3d) and (F.4b) follows from this structure, which we show in the Appendix E.

The most involved part is proving Properties (F.1) and (F.4a), corresponding to $s^* \subseteq s$ and $(\sigma \sqsubset \sigma' \lor \sigma' \sqsubset \sigma)$ respectively, which require us to establish the following helper lemmas, presented in Appendix E: By Lemma E.2, we have exactly one $\sigma_{ss}$ per $\sigma$, satisfying $\sigma_{ss} \sqsubset \sigma_{end}$, and by Lemma E.5, for any terminated $v$ and $v'$ of pushVS with $v \subset v'$, there exists $\sigma$ and $\sigma'$ such that $\sigma_{end} \sqsubset \sigma'_{end}$ and $v_{x} \not\subset \sigma_{end}$. With Lemma E.2, we can that show for every $i$ we either have $\sigma_{i} \subset \sigma_{ss} \subset \sigma_{end} \subset \sigma'_{init} \subset \sigma'_{i}$ or $\sigma'_{i} \subset \sigma'_{ss} \subset \sigma'_{end} \subset \sigma_{init} \subset \sigma_{i}$, implying Property (F.4a). With Lemma E.5, we can show for $s$ that there exists some $\sigma'$ and $\sigma''$ with $\sigma'_{end} \sqsubset \sigma''_{end}$ and $v_{x} \not\subset \sigma''_{end}$ where $v_{x} \sqsubset s_{v_{i}}$. This implies that if we have $\sigma = s^*$, then $\sigma'_{end} \sqsubset \sigma_{init}$, since some new virtual scan must have started after $\sigma'_{end}$, and by the instantiation of $[-\star]$, we have $\sigma_{ss} \xrightarrow{rf} s_{ss}$, thus for all $i$ we have $\sigma'_{end} \sqsubset \sigma_{init} \subset \sigma_{i} \subset \sigma_{b} \subset \sigma_{ss} \xrightarrow{rf} s_{ss}$, implying $\sigma \subseteq s$ and Property (F.1), the last step is further explained in Appendix E. □

5 APPLYING TO OTHER SNAPSHOT ALGORITHMS

To show the generality of our approach, we next apply it to the snapshot algorithm of Afek et al. [1993], presented as Algorithm 4. It is a single-writer, multi-scanner algorithm. At its core, scan of Algorithm 4 collects snapshots by reading the memory array $A$ twice. We denote the first read of index $i$ by $s_{k}a_{i}$, and the second by $s_{k}b_{i}$. The algorithm looks for changes in the array. If a change is detected, the reading is restarted in the next iteration. The index $k$ on the reading events identifies the iteration in which the event occurs. If no change to the array is detected between the first and second read events, then what was read is a valid snapshot, reflecting what was in the array at the moment the last index was first read in the current iteration. To ensure that changes to the array are properly recognized, each memory cell holds a version number $(A[i].ver)$ of the latest write, which is a number that is incremented each time a new value is written to the cell.

If scan detects that some index $i$ has changed twice during scanning, it can immediately terminate, returning the view of the latest write into $i (A[i].view)$. The view itself is a snapshot each write collects by calling scan before writing its value. Keeping views ensures that the snapshot methods are wait-free. scan tracks how many times index $i$ has changed by updating the local variable $moved[i]$ each time it detects a change. Each write commits its value simultaneously with incrementing the version number and updating the view.

We next use visibility relations to show that Algorithm 4 satisfies the Snapshot signature, and is thus linearizable by Lemma 2.3.

Lemma 5.1. Every execution of Algorithm 4 satisfies the Snapshot signature (Fig. 5).

Proof sketch. To prove that Algorithm 4 satisfies the Snapshot signature, we employ virtual scans as a simplification mechanism that allows us to elide write views from consideration when
establishing Snapshot. A virtual scan will denote a scan that detected no change in the array, i.e., it returned at line 21. To each abs scan, we can associate a virtual scan as follows. If the abs scan terminated because it detected no change, than it immediately is a virtual scan. If the abs scan terminated by returning a view from a write, i.e., returned at line 16, then that view itself is an abs scan which can, recursively, be associated with a virtual scan. Importantly, given abs scan $s$, the associated virtual scan $s^*$ satisfies $s^* \subseteq s$, from which we shall derive the Snapshot axioms.

Formally, a virtual scan $\sigma$ consists of rep events $\sigma_i$ and $\sigma_j$ for each $i$, where both $\sigma_i$ and $\sigma_j$ observe the same write (i.e., no change detected), and every $\sigma_i$ occurs before any $\sigma_j$:

$$\forall \sigma, i. \exists w_i. w_i.a \not\iff \sigma_i \land w_i.a \not\iff \sigma_j \quad (A.1)$$
$$\forall \sigma, i, j. \sigma_i \subseteq \sigma_j \quad (A.2)$$

If scan $s$ returned at line 21, we simply define the virtual scan $s^*$ as the set of rep events $s_k\sigma_i$ and $s_k\sigma_j$, where $k$ is the last reading iteration of $s$. If scan $s$ returned at line 16, then $s$ returns the view of some $w_i$, and we define $s^* = (w_i/\sigma)^*$. This is a well-founded recursive definition, because the ending time of the scan $w_i$ appears on the right is smaller than the ending time of $s$ on the left, and the ending times are bounded from below by 0. To see that the ending time of $w_i$ is below that of $s$, suppose otherwise. Then it must be $s \subseteq w_i.a$, because from the code of write we have that $w_i.s \subseteq w_i.a$. In particular, $s_k\sigma_i \subseteq w_i.a$ for every $k$. But, because $w_i.a$ is observed by some rep read in $s$, we also have $w_i.a \not\iff s_k\sigma_i$ for some $k$. Thus, we have an event $s_k\sigma_j$ that terminated before the event $w_i.a$ that it observes, which contradicts Property (V.1).

Next, we show that $s^* \subseteq s$. If $s$ terminated with no changes detected, this is trivial, since $s^*$ consists of the selected rep events of $s$. Otherwise, $s^* = (w_i/\sigma)^*$, for some write view $w_i$. By recursion on the definition of $[-]_s^*$ it must be $s^* = w_i/\sigma \subseteq w_i.s$, so it suffices to show $w_i.s \subseteq s$. This holds because, to reach the return at line 16, two changes of $A[i]$ must have occurred, thus there are two different writes to $A[i]$ before $w_i$ that were observed by some rep event in $s$, implying $s$ starts before $w_i$ and thus also before $w_i.s$. Additionally, since some $s_k\sigma_j$ must have read $w_i.a$, i.e., $w_i.a \not\iff s_k\sigma_j$, it must be that $w_i.s$ ends before $s$, as otherwise $s_k\sigma_j \subseteq w_i.a$ contradicts Property (V.1).

We next proceed to establish the Snapshot signature. First, we instantiate $\forall w_i$ to be the set of all writes $w_i \in W$ where $w_i.a$ is defined; these are the writes that executed their effect. Second, we
was to modularize the linearizability proofs. The work on time-stamped stack algorithm of Dodds et al. [2013] builds on this by introducing a hybrid approach, combining the visibility method together with the linearization points method. Visibility has also been applied on algorithms of sequential execution in that order. That said, we have established axioms of the signature then simply state straightforward properties of the linearization order.

In English, the scan \( s \) reads from, and also observes, \( w_i \) iff there is an appropriate rep event in the virtual scan \( s^* \) that reads from \( w_i \) at the level of rep events. We now argue that the definitions of the visibility relations satisfy the Snapshot properties. Property (V.1) holds since if \( e <^e e' \) and \( e' \subseteq e \) then \( e <^e e' \) can only hold if \( e \bullet e' \), i.e., \( w_i \lhd \sigma_i \). But then, by Property (RB.2), from \( e' \subseteq e \) we derive \( \sigma_i \supseteq w_i \), which contradicts Property (V.1) for registers. Property (S.1) holds by the structure of the algorithm. Property (S.3) holds since if \( w_i \lhd \sigma_j \) and \( w_j \lhd \sigma_i \), then \( w_i \) and \( w_j \) are observed by the same read. Thus, by Property (M.3), \( w_i = w_j \), and since each write executes \( w_i \) only once by Property (ES.2), it must be \( w_i = w_i \). Property (S.4) holds because the algorithm is assumed to be single-writer. Property (S.5) holds since each terminated write must have \( w_i \) defined, which is how we define the set \( \mathcal{W}_i \).

Next, Property (S.2) says that given \( w_i \lhd \sigma_j \), there exists no \( w_j \) such that \( w_i < w_j < s \). To prove it, we assume that \( w_j \) exists, and derive a contradiction. By assumption, we have \( w_i \lhd \sigma_j \), \( \sigma \) for \( \sigma = s^* \). From \( \sigma = s^* \subseteq s \), and \( w_i < w_j < s \), we can derive \( w_i < w_j < \sigma \), and then by Property (A.2), also \( w_i < w_j < \sigma_i \). We elide the proof of the last derivation; it is similar to Lemma 3.1 and is in Appendix F. By the definition of \( \lhd \), and Property (A.1), \( w_i \lhd \sigma \) implies \( w_i \lhd \sigma_i \). But then \( w_i \lhd \sigma_i \) occurs between \( w_i \) and \( \sigma_i \), contradicting Property (M.2) and \( w_i \lhd \sigma_i \).

Lastly, we prove Property (S.6): \( w_i, w_j \lhd s \) and \( w_j, w_j \lhd s' \) with \( w_j < w_i \) and \( w_j < w_j \) lead to contradiction. We unfold the assumptions \( w_i, w_j \lhd s \) and \( w_j, w_j \lhd s' \) into \( w_i, w_j \lhd \sigma \) and \( w_j, w_j \lhd \sigma' \), where \( \sigma = s^* \) and \( \sigma' = s^* \). By definition of \( \lhd \), and Property (A.1), from \( w_j \lhd \sigma \) we have \( w_j \lhd \sigma_j, \sigma_j \), and similarly for \( w_j \lhd \sigma' \). By Property (A.2) we have \( \sigma_j \supseteq \sigma_i \) and \( \sigma' \supseteq \sigma'j \), and by Property (RB.1), it is either \( \sigma_j \supseteq \sigma'j \) or \( \sigma'j \supseteq \sigma_i \). In the first case (the second is symmetric) we can construct \( w_j < w_j \lhd \sigma_i \), \( \sigma_i \supseteq \sigma'j \), meaning that we have \( w_j < w_j < \sigma'j \) contradicting Property (M.2) for \( w_j \lhd \sigma'j \). 

We conclude this section by observing that the inverse of Lemma 2.3 also holds; that is, Snapshot signature is actually satisfiable by every linearizable snapshot algorithm. This follows by instantiating \( < \) with the linearization order, defining \( \lhd \) to relate each read with the latest write before that read in the linearization, and defining \( \mathcal{W}_i \) as the set of all writes in the linearization. The axioms of the signature then simply state straightforward properties of the linearization order and of sequential execution in that order. That said, we have established Snapshot signature for Algorithm 4 without assuming linearizability, so that we can derive linearizability by Lemma 2.3.

6 RELATED WORK

The idea to use visibility relations for proving linearizability has first been proposed by Henzinger et al. [2013] and applied to concurrent queue algorithms. The explicit motivation of this approach was to modularize the linearizability proofs. The work on time-stamped stack algorithm of Dodds et al. [2015] builds on this by introducing a hybrid approach, combining the visibility method together with the linearization points method. Visibility has also been applied on algorithms over weak memory, where it is typically called communication order [Raad et al. 2019]. Another related approach of visibility by Emmi and Enea [2019]; Krishna et al. [2020] formulates weakened specifications for linearizability, where all operations do not have to linearly occur one after another in a simulation, but occur in partial order, allowing for certain operations to miss another.

Concerning Jayanti’s algorithms specifically, Jayanti [2005] sketched linearizability of the first algorithm by describing its linearization points. He also argued informally, based on forwarding principles, that the changes between the algorithms preserve linearizability. Delbianco et al. [2017]
gave a proof formalized in Coq of the first algorithm, by constructing a linearization order along the execution of the algorithm. The proof is specific to the implementation, making it is unclear how to lift it to the other two algorithms. Based on that development, Jacobs [2018] developed a mechanized proof, also of the first algorithm, in VeriFast using prophecy variables. Petrank and Timnat [2013] and Timnat [2015] present an algorithm based on the forwarding idea of Jayanti, which implements a set interface with insert, remove and contains operations, along with an iterator which is a generalized form of scanner. The set operations can be forwarded to the scanner, generalizing Jayanti’s forwarding which applies only to write operations. In the future, we will consider how to generalize our proof to apply to this algorithm as well. This would involve axiomatizing the set data-structure, but also what it means abstractly to be a forwarding structure, so that both Jayanti’s snapshot and the set structure of Petrank and Timnat are instances.

Afek et al. [1993] proved their algorithm linearizable by using the linearization point method. While the use of the linearization point method is fairly straightforward for this algorithm, we showed that we can reuse our snapshot axiomatization for its proof.

7 CONCLUSION AND FUTURE WORK

We presented proofs of linearizability of Jayanti’s three snapshot algorithms developed in a modular fashion, using the method of visibility relations. More concretely, the linearizability proofs are decomposed into proof modules, with many of the modules being shared between the three algorithms; they are developed once and reused three times to reduce proof complexity.

Importantly, the module interfaces are signatures consisting of relations and axioms on them that encode the key idea of “forwarding principles” underpinning Jayanti’s design. We thus show that a formalism based on visibility relations is powerful enough to mathematically capture these principles; previously, Jayanti has only presented them informally in English.

In the future, we plan to apply the visibility methods to other algorithms, and potentially also use the developed modules and signatures as guides in designing new and more efficient algorithms, much like Jayanti’s three algorithms start with the single-writer/single-scanner variant and build to the ultimately desired multi-writer/multi-scanner variant. We shall also study how the visibility method applies to non-linearizable algorithms, and how to mechanize our proof in a proof assistant.

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A SNAPSHOT SIGNATURE IMPLIES LINEARIZABILITY

Here we prove that every algorithm with histories satisfying the Snapshot signature is linearizable. This proof is based on a similar proof for queues by Chakraborty et al. [2015].

Following Definition 2.2 of linearizability, we start by choosing a visibility relation \(<\) that constructs the set \(E_c = \{ e \mid \exists e' \in \mathcal{T}(E). e < e' \}\), and then extend it to a linearization \(<\). Let \(<\) be the visibility relation obtained by restricting the visibility relation postulated by Snapshot to \(\mathcal{T}(S) \cup \cup_i \mathcal{W}_i\). Events outside of the restriction are non-terminated scans or writes that have not executed their effect yet; they do not modify the abstract state, and are hence not necessary for linearization. Additionally, by Property (S.5) we know that every terminated write is effectful. Thus going forward, we consider \(E_c = \mathcal{T}(S) \cup \cup_i \mathcal{W}_i\).

To move towards the construction of a linearization, we will first construct a partial order \(\prec_w\) over all write events that contains all logically implied orderings. We will in steps expand this order, until we have a total order.

**Definition A.1.** Let \(w_i \xrightarrow{wrDiff} w_j\) be defined if \(i \neq j\) and there exists a write \(w_j\) and a scan \(s\) such that \(w_i, w_j \xrightarrow{rf} s\) and \(w_j \prec w_j\). Let \(\prec_w = (\prec \cup \xrightarrow{wrDiff})^+\).

**Lemma A.2.** The relation \(\prec_w\) is a strict partial order.

**Proof.** For a relation to be a strict partial order, it needs to be transitive and irreflexive. The relation \(\prec_w\) is transitive by definition since it is a transitive closure. We prove irreflexivity by proving acyclicity of \(\prec_{w1} = \prec \cup \xrightarrow{wrDiff}\), i.e., we prove there exists no cycle of the following form:

\[w_1 \prec_{w1} w_2 \prec_{w1} \cdots \prec_{w1} w_n \prec_{w1} w_1\]

We show this by showing that there exists no cycle of size \(n = 1\) and that we can shrink any cycle of size \(n \geq 2\), meaning that we can recursively shrink a cycle until it is of size \(n = 1\). For a cycle with \(n = 1\), we have \(w_1 \prec_{w1} w_1\), which can only be the case \(w_1 \prec w_1\), since two writes joined by \(\xrightarrow{wrDiff}\) must come from distinct memory cells. However, we cannot have \(w_1 \prec w_1\) either, since it contradicts \(\prec\) being irreflexive (Lemma 2.1). We show that we can shrink any cycle of \(n \geq 2\) by doing a case analysis on the first two instances of \(\prec_{w1}\) in the cycle:

1. We have \(w_1 \prec w_2 \prec w_3 \prec_{w1} w_1\). By transitivity of \(\prec\) we can shrink the cycle to the smaller cycle \(w_1 \prec w_2 \prec w_3 \prec_{w1} w_1\).
2. We have \(w_1 \xrightarrow{wrDiff} w_2 \prec w_3 \prec_{w1} w_1\), i.e., \(w_1 = w_i\) and \(w_2 = w_j\) with \(i \neq j\) and there exists \(w_j\) such that \(w_1, w_j \xrightarrow{rf} s\) and \(w_j \prec w_j\). Let \(w_3 = w_k\) where \(k\) can either be equal to \(i\) or \(j\) or some other distinct index. Since each scan is terminated and thus observes the entire memory by Property (S.1), there must exist some \(w_k\) such that \(w_j \xrightarrow{rf} s\). By comparing \(w_k\) and \(w_k\) with Property (S.4), we have either \(w_k \prec w_k\) or \(w_k \prec w_k\). If we have \(w_k \prec w_k\) then we also have \(w_j \prec s\) by \(w_j \prec w_k\) and \(w_k \xrightarrow{rf} s\), which implies that \(w_j\) happens between \(w_j\) and \(s\), contradicting Property (S.2). Thus we must have \(w_k \prec w_k\), which lets us construct \(w_1 \xrightarrow{wrDiff} w_3 \prec_{w1} w_1\), giving us a shorter cycle.
3. We have \(w_1 \xrightarrow{wrDiff} w_2 \xrightarrow{wrDiff} w_3 \prec_{w1} w_1\), i.e., \(w_1 = w_i\) and \(w_2 = w_j\) and \(w_3 = w_k\) with \(w_i, w_j \xrightarrow{rf} s\) and \(w_j, w_k \xrightarrow{rf} s'\) and \(w_j \prec w_j\) and \(w_k \prec w_k\) and \(w_3 \prec_{w1} w_1\) where \(i \neq j\) and \(j \neq k\). Similar to the previous case, by Property (S.1), there must exist a \(w_k\) such that \(w_j \xrightarrow{rf} s\). By Property (S.6) with \(w_j, w_k \xrightarrow{rf} s\) and \(w_j, w_k \xrightarrow{rf} s'\) and \(w_j \prec w_j\), it must follow that \(w_k \prec w_k\), which by Property (S.4) implies \(w_k \xrightarrow{rf} w_k\) giving us \(w_k \prec w_k\) by \(w_k \xrightarrow{rf} w_k\). This lets us construct \(w_1 \xrightarrow{wrDiff} w_3 \prec_{w1} w_1\) giving us a shorter cycle.
4. We have \(w_1 \prec w_2 \xrightarrow{wrDiff} w_3 \prec_{w1} w_1\). For this case we simply rotate the cycle to the first instance to the end of the cycle, i.e., \(w_2 \xrightarrow{wrDiff} w_3 \prec_{w1} w_1 \prec w_2\), and case match on the next two \(\prec_{w1}\) instances to create a scenario where we either use case (2) or case (3).
By Zorn’s Lemma, we can arbitrarily choose a total order $<_w$ extending $<_w$, for which we define the concept of maximal candidate, which we use to construct the complete total order.

**Definition A.3.** An event $e$ is called a maximal candidate of $<_w$ when $e$ is $<_w$-maximal (i.e., there exists no greater event in $<_w$, which may not necessarily be unique) and one of the following holds:

1. $e = w_i$ is a write with no scans observing it and it is the greatest write in $<_w$,
2. $e = s$ is a scan with for all $i$, there exists $w_i$ such that $w_i \rightarrow s$, and $w_i$ is the greatest event of order $<_w$ in $\mathcal{H}_i$.

**Lemma A.4.** For any non-empty history satisfying the Snapshot signature, there exists a maximal candidate of $<_w$.

**Proof.** Since $<_w$ is a total order over a finite set, there exists a unique greatest write $w_i$, consider the following cases:

1. If nothing reads $w_i$ and $w_i$ is $<_w$-maximal, then $w_i$ is trivially a maximal candidate.
2. If nothing reads $w_i$ and $w_i$ is not $<_w$-maximal, then there exists an event $e$ greater than $w_i$, such that $w_i < e$. The event cannot be a write event $e = w_j'$, since that would imply $w_i < w_j'$, contradicting $w_i$ being $<_w$-maximal. Thus it must be a scan event $e = s$, by Property (S.1) there exists $w_j'$ such that $w_j' \rightarrow s$ with $w_j' \neq w_i$, since nothing reads $w_i$. Since we have $w_j' \rightarrow s$ and $w_i < s$, it must be the case that $w_i$ does not happen in-between $w_j'$ and $s$ as per Property (S.2), i.e., we have $w_i' \neq w_j$, thus by Property (S.4) we have $w_i < w_i'$. With $w_i$ being $<_w$-maximal, we also have $w_i' < w_i$, however this contradicts $w_i < w_i'$.
3. If there exists $s$ such that $w_i \rightarrow s$, with $s$ being $<_w$-maximal, then we show that $s$ is a maximal candidate For $s$ to be a maximal candidate, we need for all $j$, there exists $w_j$ such that $w_j < s$, $w_j$ is $<_w$-maximal in $\mathcal{H}_j$. By Property (S.1), we know that there exists some $w_j$ such that $w_j \rightarrow s$. Suppose for each $j \neq i$, $w_j$ is not $<_w$-maximal in $\mathcal{H}_j$, then there exists $w_j'$ such that $w_j < w_j'$. This implies $w_j < w_j'$ by Property (S.4) and since $<_w$ extends $<_w$, which by $w_i, w_j \rightarrow s$ and $w_j < w_j'$, it implies $w_i < w_j'$, however this contradicts maximality of $w_i$. Thus each $w_j$ is $<_w$-maximal in $\mathcal{H}_j$, meaning $s$ is a maximal candidate.
4. If there exists $s$ such that $w_i \rightarrow s$, with $s$ not being $<_w$-maximal, then there must exist some other event $e$ such that $s < e$. If there exists multiple scans observing $w_i$ we let $s$ be the $<_w$-greatest among them. If $e$ is a write event $e = w_j'$ then we have $w_i < w_j'$ since $w_i \rightarrow s < w_j'$, which contradicts $w_i$ being the greatest write in $<_w$. If $e$ is a scan $e = s'$ then we either have $w_i \rightarrow s'$ or $w_i \rightarrow s'$ for some write $w_j \neq w_i$, however the former contradicts our selection of $s$ since both $s$ and $s'$ observe $w_i$, but $s < s'$ contradicts that $s$ be the $<_w$-greatest among scans observing $w_i$. We thus have that $w_j \rightarrow s'$ and since we also have $w_i < s'$ by $w_i \rightarrow s < s'$, if $w_j$ is in-between $w_j'$ and $s'$ we contradict Property (S.2), thus we have $w_i' \neq w_j$ which by Property (S.4) we must have $w_i < w_i'$, which contradicts that $w_i$ is $<_w$-maximal. □

To construct the total order, we recursively select a maximal candidate, which we make the last event of the total order and remove from consecutive steps.

**Theorem A.5.** Every history satisfying our snapshot specifications is linearizable, i.e., there exists a total order $<$ over each event which forms a legal sequential snapshot behavior and extends $<_w$ and $<$. 

**Proof.** By induction on the number of events in the history. The minimal history, the one that only consists of the initial writes, is clearly linearizable, no matter what order $<_w$ we have.

Our inductive hypothesis is that for some history $E \subseteq E_\mathcal{C}$ satisfying the axioms under $<_w$, there exists a total order $<$ over $E$ that forms a legal sequential snapshot behavior and extends $<_w$ and $<$ when only considering events in $E$. For the inductive step, we need to prove for a there exists
such a total order $<'$ of history $E'$, which extends $E$ by the single event $e$, with $e$ being an arbitrary maximal candidate of $<_w$ under history $E'$. By Lemma A.4, since $E'$ is non-empty, there must exist a maximal candidate. Consider the two following cases for $e$.

- $e = w_i$ is a write. By the definition of being a maximal candidate, we know that there exists no scan in $E'$ observing $w_i$. If we remove $w_i$ from $E'$ then the new history $E$ still satisfies the snapshot properties since it is safe to remove any event that is not read visible. By the inductive hypothesis we have a total order $<$ over $E$ that forms a legal snapshot behavior and extends $<_w$ and $\prec$. To construct the total order $<_'$ over $E'$, we extend $<$ with write $w_i$ being the last event in the order. Since any write event added to the end of $<_'$ cannot influence the result of previous scan events, it means that since $<$ forms a legal snapshot behavior, so must $<_'$.

- $e = s$ is a scan. By the definition of being a maximal candidate, we know that all the writes $s$ observe are maximal under $<_w$ in their respective memory cell. If we remove $s$ from $E'$ then the new history $E$ still satisfies the snapshot properties since no other events depend on scans. By the inductive hypothesis we have a total order $<$ over $E$ that forms legal snapshot behavior and extends $<_w$ and $\prec$. To construct the total order $<_'$ over $E'$, we extend $<$ with scan $s$ being the last event in the order. The $<_'$ order forms a legal snapshot behavior, since any write $w_i$ that $s$ observes is maximal in $<_w$, it is impossible for another write to be in-between $w_i$ and $s$, thus since $<$ forms a legal snapshot behavior, so does $<_'$.

\[\Box\]

**B FORWARDING SIGNATURE IMPLIES SNAPSHOT SIGNATURE**

We here give the full proof that any history satisfying the Forwarding signature (Fig. 7) also satisfies the Snapshot signature (Fig. 5). We start by determining the instantiations of $\mathbb{W}_i$ and $\mathbb{WF}_i$ and $<$ for the Snapshot signature. Let $w_i \in \mathbb{W}_i$ iff $w_i$ has executed $w_{ij}a$, since the effect of $w_i$ is only observable after this execution. We define a helper notion of visibility over scans $\mathbb{vs}$, which we use to define $\prec$. The highlighted relations $\mathbb{WF}_i$, and $\mathbb{WF}_i$, originate from the Forwarding signature.

\[
\begin{align*}
\text{\textbf{R.1}} & \quad w_i \xrightarrow{\mathbb{WF}_i} s \Leftrightarrow w_i \xrightarrow{\mathbb{WF}_i} s^* \\
\text{\textbf{R.2}} & \quad s \xrightarrow{\mathbb{vs}} s' \Leftrightarrow s^* \sqsubseteq s^* \\
\text{\textbf{R.3}} & \quad w_i \xleftarrow{\mathbb{MF}_i} \Leftrightarrow w_i \xleftarrow{\mathbb{WF}_i} \\
\text{\textbf{R.4}} & \quad w_i \xrightarrow{\mathbb{MF}_i} s \Leftrightarrow w_i \xrightarrow{\mathbb{WF}_i} s
\end{align*}
\]

We prove the Snapshot properties by unfolding the definition of $\mathbb{WF}_i$ in the properties so that we can apply the Forwarding properties.

**Property (S.1).** For every terminated scan $s$ and index $i$, there exists some write $w_i$ such that $w_i \xrightarrow{\mathbb{WF}_i} s$ and $w_i.in = s.out[i]$.\[\text{\textbf{Proof.}}\] By Property (F.2), there exists $w_i$ such as $w_i \xrightarrow{\mathbb{WF}_i} s^*$ and $w_i.in = s.out[i]$ for $s$, where the latter immediately corresponds to one of our goals and the former corresponds to the remaining goal by the instantiation of $\mathbb{WF}_i$, giving us $w_i \xrightarrow{\mathbb{WF}_i} s = w_i \xrightarrow{\mathbb{WF}_i} s$.\[\Box\]

**Property (S.5).** Every terminated write is effectful, i.e., $\mathcal{T}(W_i) \subseteq \mathbb{W}_i$.\[\text{\textbf{Proof.}}\] Since every terminated event must have executed all their rep events, it follows that a terminated write $w_i$ must have executed $w_i.a$, which matches the instantiation of $\mathbb{W}_i$.\[\Box\]

**Property (S.4).** For two distinct writes $w_i$ and $w'_i$, we have either $w_i \prec w'_i$ or $w'_i \prec w_i$.\[\text{\textbf{Proof.}}\] For each write, we either have $w_i.a \prec w'_i.a$ or $w'_i.a \prec w_i.a$ by Property (M.4). By the definition of $\mathbb{WF}_i$, if we have $w_i.a \prec w'_i.a$ then we have $w_i \xrightarrow{\mathbb{WF}_i} w'_i$ (and thus also $w_i \prec w'_i$) and if we have $w'_i.a \prec w_i.a$ then we have $w'_i \xrightarrow{\mathbb{WF}_i} w_i$ (and thus also $w'_i \prec w_i$).\[\Box\]
**Property (V.1).** If \( e \not<+ e' \) then we cannot have \( e' \sqsubseteq e \).

**Proof.** Assume we have \( e' \sqsubseteq e \), we derive a contradiction. Consider first that we have \( e = e' \) such that we have \( e \not<+ e' \). There is no \( w_j \not\sqsubseteq s \) in \( e \not<+ e \) since if there is, there is no way to relate from a scan back to a write, i.e., there is no observation from scan to write. Thus we must have either \( e \not\sqsubseteq e \) or \( e \not\sqsubseteq e' \), which we cannot have by total ordering of \( w_j a \) events by Property (M.4) or by virtual scan total order by Property (F.4a). Thus \( e \) and \( e' \) cannot be equal.

Assuming instead we have \( e' \sqsubseteq e \), we derive a contradiction. By Lemma 3.1 from \( e \not<+ e' \) we have \( e_r \not<_r e_r' \) for some \( e_r \sqsubseteq e \) and \( e_r' \sqsubseteq e' \). By Property (RB.2) we can derive \( e'_r \sqsubseteq e_r \) from \( e' \sqsubseteq e \), we can construct the cycle \( e_r' \not<_r e_r' \) which contradicts irreflexivity of \( <_r \) (Lemma 2.1). \( \square \)

By Lemma 2.1 and Virtual Property (V.1), we have that \( <_r \) is irreflexive, and thus it is a strict partial order. We will use this to prove following properties.

**Property (S.3).** If \( w_j \not\sqsubseteq s \) and \( w'_j \not\sqsubseteq s \) then we must have \( w_i = w'_i \).

**Proof.** Let \( s^* = \sigma \), by unfolding \( w_j \not\sqsubseteq s \) we have \( w_j \not\sqsubseteq s \), \( \sigma \), and dually for \( w'_j \not\sqsubseteq s \). By case splitting \( w_j \not\sqsubseteq s \) we have \( w_j \not\sqsubseteq s \) and \( w'_j \not\sqsubseteq s \). Remember that \( w_j \not\sqsubseteq s \) splits into either \( w_j a \not\sqsubseteq \sigma_j = \sigma_l \) and \( \sigma^* = \sigma \), or \( w_j \not\sqsubseteq \sigma_j = \sigma_l \), or \( w_j \not\sqsubseteq \sigma_j = \sigma_i \) and \( \sigma^* = \sigma \), or similarly for \( w'_j \not\sqsubseteq s \).

- If \( w_j \not\sqsubseteq s \) and \( w'_j \not\sqsubseteq s \) splits into different cases, e.g., we have \( w_l \not\sqsubseteq s \) and \( \sigma_l \not\sqsubseteq \sigma_l \), then by Property (F.5b) we derive a contradiction.
- If both split into the first case, i.e., we have \( w_j a \not\sqsubseteq \sigma_j = \sigma_l \) and \( w_j a \not\sqsubseteq \sigma_j = \sigma_l \), then by Property (M.3) we have \( w_j a = w'_j a \). Since each abstract write has a distinct representative write, we must have \( w_j = w'_j \).
- If both split into the second case, i.e., we have \( w_l \not\sqsubseteq s \) and \( w'_l \not\sqsubseteq s \) then by Property (F.5a) we have \( w_l = w'_l \).

**Lemma B.1.** If \( w_i < s \) then \( w_i \not<_r s \).

**Proof.** By viewing \( w_i < s \) as a chain of \( <_r \), we have \( w_i <_1 s, \ldots <_1 s \). By induction on the length of the chain. In our base case we have \( w_i <_1 s \), which splits into cases \( w_i \sqsupset s \) and \( w_i < s \), which can only be \( w_i \not\sqsubseteq s \). For \( w_i \sqsubseteq s \), by Properties (RB.2) and (F.1), we have \( w_i \sqsubseteq s^* \). For \( w_i \not\sqsubseteq s \), we have \( w_i \not\sqsubseteq s^* \), \( s^* \) directly from the definition of \( \not\sqsubseteq \). The inductive step is similar, with the exception that we need to also consider for \( w_i <_1 \ldots <_1 s' <_1 s \) the \( s' \not\sqsubseteq s \) case, which by the definition of \( \not\sqsubseteq \) gives us \( s^* \sqsubseteq s^* \), which implies \( s^* \not<_r s^* \), allowing us to construct \( w_i < s^* \).

**Property (S.2).** If \( w_j \not\sqsubseteq s \) then there does not exist a write \( w'_j \) such that \( w_i < w'_j < s \) holds.

**Proof.** Assume we have such a write \( w'_j \), we derive a contradiction. Let \( \sigma = s^* \), we unfold \( w_j \not\sqsubseteq s \) into \( w_j \not\sqsubseteq s \), \( \sigma \). By Lemma B.1 we have \( w_j < s \) from \( w'_j < s \). We first show that there exists a write \( w''_j \) with \( w_j < w''_j \) such that either \( w''_j a < \sigma_j \) or \( w''_j < \sigma_j \) hold. We split \( w_j < \sigma_j \) into multiple cases, multiple times, as shown in Fig. 12. We justify each of the implications:

- Any implication from something with \( <_r \) into either \( <_1 \sqsubseteq \) or \( \not<_1 \) simply holds by definition.
- The implication \( w'_j \not\sqsubseteq s \) into \( w_j \not\sqsubseteq s \) holds since there has to exist a write \( w_j \) and a scan \( \sigma' \) such that \( w'_j < \sigma_j \) and \( \sigma' < w_j \). By Property (F.4a), we either have \( \sigma \not<_r \sigma' \) or \( \sigma' \not<_r \sigma \), where in the former case we have \( \sigma' < w_j \) and \( \sigma \not<_r \sigma' \), contradicting irreflexivity of \( <_r \), thus we must have the latter. Thus we have \( w'_j < \sigma_j \).
- All the implications going into \( w'_j < \sigma_j \) holds by Lemma 3.1 and Property (RB.2).

We now consider each of the possible final cases of Fig. 12:
Fig. 12. Directed acyclic graph representing the proof of splitting \( w'_i \prec \sigma \) into four final (boxed) cases. Forks in the graph with \( \rightarrow \) arrows represents case splitting, splitting either \( \triangleleft_1 \) or \( \triangleleft \). Implication arrows (\( \Rightarrow \)) represents an implication.

- If we have \( w'_i \overset{wr}{\triangleright} \sigma \), then we must have some \( w''_i \) such that \( w'_i \overset{wr}{\triangleright} w''_i \prec \sigma \). We recursively consider the cases of Fig. 12 with \( w_i \prec w''_i \prec \sigma \). We cannot have an infinite recursive chain of this case, since the set of writes is finite and thus, the only way to have such an infinite chain is by having a cycle in \( \overset{wr}{\triangleright} \), which is impossible since such a cycle implies we also have a cycle of rep events.
- If we have \( w'_i \overset{rf}{\triangleright} \sigma \), then we contradict the result of proving Property (S.3) and \( w_i \overset{rf}{\triangleright} \sigma \) since \( w_i \) and \( w'_i \) cannot be equal by \( w_i \prec w'_i \).

Thus the only remaining cases are \( w''_i a \prec \sigma i \) and \( w''_i \overset{\triangleleft}{\triangleleft} \sigma \). We case split \( w_i \overset{rf}{\triangleright} \sigma \) and consider all possible combined cases:

- If \( w_i a \overset{rf}{\triangleright} \sigma i \) and \( w'_i a \prec \sigma i \), then \( w''_i a \) happens in-between \( w_i a \) and \( \sigma i \), thus we contradict Property (M.2).
- If \( w_i a \overset{rf}{\triangleleft} \sigma i \) and \( \sigma i \overset{rf}{\triangleleft} \sigma b j \) and \( w''_i \overset{\triangleleft}{\triangleleft} \sigma \), then from Property (F.6) we have \( w''_i a \prec \sigma a i \), implying \( w''_i a \) happened in-between \( w_i a \) and \( \sigma a i \), thus we contradict Property (M.2).
- If \( w_i \overset{fwd}{\triangleright} \sigma \) and \( w''_i a \prec \sigma i \), then Property (F.7a) contradicts \( w_i \prec w''_i \).
- If \( w_i \overset{fwd}{\triangleleft} \sigma \) and \( w''_i \overset{\triangleleft}{\triangleleft} \sigma \), then Property (F.7b) contradicts \( w_i \prec w''_i \).

**Property (S.6).** If \( w_i, w_j \overset{rf}{\triangleleft} s \) and \( w'_i, w'_j \overset{rf}{\triangleleft} s' \) with \( w_i \prec w'_i \) then we cannot have \( w'_i \prec w_j \).

**Proof.** Assume we have \( w'_i \prec w_j \), we derive a contradiction. Let \( \sigma = s^* \) and \( \sigma' = s'^* \), we unfold \( w_i, w_j \overset{rf}{\triangleleft} s \) into \( w_i, w'_j \overset{rf}{\triangleleft} s^* \) and similarly for \( w'_i, w'_j \overset{rf}{\triangleleft} s' \). By Property (F.4a), we either have \( \sigma = \sigma' \) or \( \sigma \subset \sigma' \) or \( \sigma' \subset \sigma \). If we have \( \sigma = \sigma' \), we contradict the result of proving Property (S.3), since we would have \( w_i \overset{rf}{\triangleleft} \sigma \) and \( w_j \overset{rf}{\triangleleft} \sigma \) and \( w_i \overset{rf}{\triangleleft} \sigma \) and \( w_j \overset{rf}{\triangleleft} \sigma \) and \( w'_i \overset{rf}{\triangleleft} \sigma \) and \( w'_j \overset{rf}{\triangleleft} \sigma \), but \( w_i \overset{rf}{\triangleleft} \sigma \) and \( w'_i \overset{rf}{\triangleleft} \sigma \) cannot be equal since we have \( w_i \prec w'_i \). If we have \( \sigma \subset \sigma' \), then we have \( w'_i \prec w_j \overset{rf}{\triangleright} \sigma \subset \sigma' \), contradicting the result of proving Property (S.2) by \( w_j \) occurring in between \( w'_i \overset{rf}{\triangleright} \sigma' \). Similarly, if we have \( \sigma' \subset \sigma \), then we have \( w_i \prec w'_i \overset{rf}{\triangleright} \sigma' \subset \sigma \), contradicting the result of proving Property (S.2) by \( w'_i \) occurring in between \( w_i \overset{rf}{\triangleright} \sigma \). \( \square \)

**C Algorithm 1 Satisfies the Forwarding Signature**

We here present the complete proof for Jayanti’s single-writer/single-scanner algorithm (Algorithm 1) satisfying the Forwarding signature (Fig. 7).

Since this algorithm is single-scanner, we instantiate the set of virtual scans to be equal to the set of abs scans, \( \Sigma = S \), and let \([-1]^* \) be the identity mapping, \( s^* = s \). We let \( A \) correspond to the
array of the same name and $B$ for each virtual scan will only map to $B$. The rep events of the event signatures for write and vscan (Fig. 6) corresponds to the rep events of the algorithm of the same name, i.e., $w_i a$ and $sr_j$, $sa_i$ and $sh_i$. We instantiate forwarding visibility to the following:

$$w_i \xrightarrow{fwd} s \iff w_i b \xrightarrow{fwd} sh_i$$

I.e., a write $w_i$ was forwarded to scan $s$ if $sh_i$ read the value written by $w_i b$.

By $[-]^*$ being instantiated as the identity mapping, Property (F.1) holds, since for all $e$, $e \subseteq e$. By the algorithm being single-scanner implies that virtual scan total order Property (F.4a) holds. The structure of the algorithm and the instantiation of $fwd$, implies that Properties (F.2) and (F.4b) hold. Additionally, since the algorithm only writes to $A$ in writers and writes $\bot$ to $B$ in scanners, Properties (F.3a) and (F.3b). This leaves Properties (F.5a) to (F.7b) to prove.

**Property (F.5a).** If we have $w_i \xrightarrow{fwd} s$ and $w_j \xrightarrow{fwd} s$ then $w_i$ and $w_j$ must be equal.

Proof. By our instantiation we have $w_i b \xrightarrow{fwd} sb_i$ and $w_j b \xrightarrow{fwd} sb_i$, thus by Property (M.3) we have $w_i b = w_j b$. Since all writes only performs their operations at most once, this means that the writes must be the same, i.e., $w_i = w_j^\prime$. □

**Property (F.5b).** If $sh_i$ terminated without $sr_j \xrightarrow{fwd} sb_i$ then there exists $w_i$ such that $w_i \xrightarrow{fwd} s$.

Proof. Since we do not have $sr_j \xrightarrow{fwd} sb_i$, it is the case that $sb_i$ could not have observed any other $s' r_j$ since such $s'$ must have occurred before $s$ by single-scanner, implying $s' r_j \subseteq sr_j$, meaning we would contradict Property (M.2) if we had $s' r_j \xrightarrow{fwd} sb_i$ by $sr_j$ occurring in between them. Thus, $sb_i$ must have observed some $w_i b$, i.e., $w_i b \xrightarrow{fwd} sb_i$, implying our goal $w_i \xrightarrow{fwd} s$. □

**Property (F.5c).** If we have $w_i \xrightarrow{fwd} s$ then we have $w_i a \prec sb_i$ without $sr_j \xrightarrow{fwd} sb_i$.

Proof. Since $w_i \xrightarrow{fwd} s$ is instantiated as $w_i b \xrightarrow{fwd} sb_i$, by $w_i a \equiv w_i b \xrightarrow{fwd} sb_i$ we have $w_i a \prec sb_i$. For proving absence of $sr_j \xrightarrow{fwd} sb_i$, assume we have it, we derive a contradiction. By $w_i \xrightarrow{fwd} s$ we have $w_i b \xrightarrow{fwd} sb_i$, which by Property (M.3) we have $w_i b = sr_j$, however that is impossible, since $w_i b$ is a rep event of $w_i$, while $sr_j$ is a rep event of $s$, thus a contradiction. □

Note that we now satisfies Properties (F.4b) and (F.5c), meaning we satisfy the preconditions to use Lemma 3.1. Before we prove Properties (F.6), (F.7a) and (F.7b), we prove some helper lemmas.

**Lemma C.1.** For any execution, we have $s_{on} \xrightarrow{rf} w_i a$ iff we have $s_{on} \prec w_i a$ and $s_{off} \not\prec w_i a$.

Proof.

(⇒) $s_{on} \xrightarrow{rf} w_i a$ trivially implies $s_{on} \prec w_i a$, and if we were to have $s_{off} \prec w_i a$ we contradict Property (M.2) by $s_{off}$ occurring in between $s_{on} \xrightarrow{rf} w_i a$.

(⇐) By the fact that only scans change $X$ and there is a single-scanner constraint, we know that any write to $X$ other than $s_{on}$ and $s_{off}$ either happens before or after the scan $s$. Since $w_i a$ has to observe something when it finishes, it must observe $s_{on}$, since if $e \xrightarrow{rf} w_i a$ for some write event $e \neq s_{on}$, then by Property (M.4) we either have $e \prec s_{on}$, which contradicts Property (M.2) by $s_{on}$ occurring in between $e \xrightarrow{rf} w_i a$; or $s_{on} \prec e$, which by the algorithm structure $e = s_{off}$ or some later write, however this contradicts $s_{off} \not\prec w_i a$. □

**Lemma C.2.** If $w'_i \equiv s \xrightarrow{rf} s$ then $s_{off} \not\prec w_i a$ and if $w'_i b$ was executed, then $w'_i b \prec sh_i$.

Proof. For $w'_i \equiv s \xrightarrow{rf} s$, there must exist some $e$ and $w_j$ such that $w'_i \equiv e \equiv w_j \xrightarrow{rf} s$. Each of these events are populated except for $s$, since each of the events are to the left of some $\prec$. For $w'_i \equiv e \equiv w_j$, let $e_r$ be a rep event executed by $e$, then by Property (RB.2) and Lemma 3.1 we have $w'_i \equiv e_r \equiv w_j a$. We case split $w_j \xrightarrow{rf} s$ into either $w_j a \xrightarrow{rf} sa_j$ or $w_j b \xrightarrow{rf} sh_j$. 
• If we have \( w_j a \xrightarrow{t} sa_j \) then we have the chain
\[
 w_j' x \sqsubset w_j' b \overset{\text{true}}{\sqsubset} w_j a \xrightarrow{t} sa_j \sqsubset s_{off} \sqsubset sb_i
\]
which trivially implies \( w_j' b \sqsubset sb_i \) and \( s_{off} \not\sqsubset w_j' x \).
• If we have \( w_j b \xrightarrow{t} sb_j \) then since \( w_j b \) wrote, we know that \( w_j x \) read \( \text{true} \), i.e., \( s'_{on} \xrightarrow{t} x \) \( w_j x \) for some \( s' \). By single-scanner, we have either \( s \sqsubset s' \) or \( s' \sqsubset s \), which in the former case we can form the cycle
\[
 sb_j \sqsubset s'_{on} \xrightarrow{t} w_j x \sqsubset w_j b \xrightarrow{t} sb_j
\]
contradicting irreflexivity of \( \triangleleft \), thus we must have \( s' \not\sqsubset s \). By Lemma C.1 we have \( s'_{off} \not\sqsubset w_j x \), which by \( s' \sqsubset s \) implies \( s_{off} \not\sqsubset w_j x \). By Property (RB.1) over \( w_j' b \sqsubset e_r \) and \( s_{off} \sqsubset sb_i \) we have \( w_j' b \sqsubset sb_i \) since if we were to have \( s_{off} \sqsubset e_r \), then we contradict \( s_{off} \not\sqsubset w_j x \) by \( s_{off} \sqsubset e_r \sqsubset w_j x \).

For the \( s_{off} \not\sqsubset w_j' x \) goal, assume we have \( s_{off} \not\sqsubset w_j' x \), then we derive a contradiction by having \( s_{off} \not\sqsubset w_j' x \sqsubset w_j x \) which contradicts \( s_{off} \not\sqsubset w_j x \).

\[\Box\]

**Property (F.6).** If \( w_j' \sqsubset \xrightarrow{t} s \) and \( sr_i \xrightarrow{t} sb_j \) then \( w_j a \not\sqsubset sa_i \).

**Proof.** By Lemma C.2 we have \( s_{off} \not\sqsubset w_j' b \) and \( w_j' b \not\sqsubset sb_j \) if \( w_j' b \) was executed. Consider if we have \( s_{on} \not\sqsubset w_j' x \), then by Lemma C.1 we have \( s_{on} \xrightarrow{t} x \) \( w_j' x \) which in turn implies that \( w_j' b \) must have executed, however that contradicts Property (M.2) by \( w_j' b \) occurring in between \( sr_i \xrightarrow{t} sb_j \), thus we must have \( s_{on} \not\sqsubset w_j' x \).

By Property (RB.1) over \( w_j' a \sqsubset w_j' x \) and \( s_{on} \not\sqsubset sa_i \), since \( s_{on} \not\sqsubset w_j' x \) contradicts \( s_{on} \sqsubset w_j' x \), we must have \( w_j' a \sqsubset sa_i \), which is our goal, thus we are done.

**Properties (F.7a) AND (F.7b).** If \( w_i \xrightarrow{\text{wd}} s \) and either \( w_j' \sqsubset \xrightarrow{t} s \) or \( w_i a \not\sqsubset sr_i \) then \( w_j \not\sqsubset w_j' \).

**Proof.** By our instantiation we have \( w_j b \xrightarrow{t} sb_j \) from \( w_i \xrightarrow{\text{wd}} s \). Assuming we have \( w_i a \not\sqsubset w_j' \) (thus we also have \( w_i a \not\sqsubset w_j' a \) by Lemma 3.1), we derive a contradiction.

Let us start with the case where we have \( w_j' a \not\sqsubset sr_i \). By the single-writer constraint we have \( w_j b \not\sqsubset w_j' a \not\sqsubset sr_i \sqsubset sb_j \) which contradicts Property (M.2) by \( sr_i \) occurring in between \( w_i \xrightarrow{t} sb_j \), thus this case is not possible. For the next case we assume \( w_j' a \not\sqsubset sr_i \).

For the \( w_j' \sqsubset \xrightarrow{t} s \) case, by Lemma C.2 we have \( s_{off} \not\sqsubset w_j' b \) and \( w_j' b \not\sqsubset sb_j \) if \( w_j' b \) was executed. By Property (RB.1) over \( w_j' a \sqsubset w_j' x \) and \( s_{on} \sqsubset sr_i \), since we have \( w_j' a \not\sqsubset sr_i \), contradicting \( w_j' a \sqsubset sr_i \) we must have \( s_{on} \sqsubset w_j' x \). By Lemma C.1 we thus have \( s_{on} \xrightarrow{t} w_j' x \), which means that \( w_j' x \) reads \( \text{true} \) and thus \( w_j' b \) is executed. By the single-writer constraint we have \( w_j b \not\sqsubset w_j' b \not\sqsubset sb_i \) which contradicts Property (M.2) by \( w_j' b \) occurring in between \( w_i \xrightarrow{t} sb_j \).

\[\Box\]

### D MULTI-WRITER FORWARDING SATISFIES THE FORWARDING SIGNATURE

We here give the complete proof that any history satisfying the MWForwarding signature (Fig. 11) also satisfies the Forwarding signature (Fig. 7). We instantiate \( \Sigma \) and \([-]* \) of Forwarding to be the same as the corresponding instantiations of MWForwarding. Properties (F.1), (F.2), (F.3a), (F.3b) and (F.4a) are shared between the signatures, therefore they trivially hold.

**Property (F.4b).** Every virtual scan satisfies \( \sigma_{r_i} \sqsubset \sigma_{a_i} \sqsubset \sigma_{b_i} \).

**Proof.** By Property (F+.4b), we have
\[
\sigma_{r_i} \sqsubset \sigma_{on} \xrightarrow{t} \sigma_{a_i} \sqsubset \sigma_{off} \xrightarrow{t} \sigma_{off} \sqsubset \sigma_{on} \xrightarrow{t} \sigma_{off} \sqsubset \sigma_{b_i}
\]
by Property (RB.1) over \( \sigma_{r_i} \sqsubset \sigma_{on} \) and \( \sigma_{on} \sqsubset \sigma_{a_i} \), we either have \( \sigma_{r_i} \sqsubset \sigma_{a_i} \) or \( \sigma_{on} \sqsubset \sigma_{on} \), however the latter contradicts \( \sigma_{on} \xrightarrow{t} \sigma_{on} \), thus we must have \( \sigma_{r_i} \sqsubset \sigma_{a_i} \). Similarly, by Property (RB.1) over
\[\sigma_u \sqsubset \sigma_{off}\text{ and }\sigma_{off} \sqsubset \sigma_b,\text{ we either have }\sigma_u \sqsubset \sigma_b\text{ or }\sigma_{off} \sqsubset \sigma_{off},\text{ where the latter contradicts }\sigma_{off} \not\sqsubseteq \sigma_{off},\text{ thus we must have }\sigma_u \sqsubset \sigma_b.\]

**Property (F.5a).** If we have \(w_i \xrightarrow{\text{fwd}} \sigma\) and \(w'_i \xrightarrow{\text{fwd}} \sigma\) then \(w_i\) and \(w'_i\) must be equal.

**Proof.** By the instantiation we have \(w_i a \xrightarrow{\text{fwd}} f a\) and \(f b \xrightarrow{\text{fwd}} \sigma b\) and \(w'_i a \xrightarrow{\text{fwd}} f' a\) and \(f' b \xrightarrow{\text{fwd}} \sigma b\) from \(w_i \xrightarrow{\text{fwd}} \sigma\) and \(w'_i \xrightarrow{\text{fwd}} \sigma\) respectively. By Property (M.3) over \(f b \xrightarrow{\text{fwd}} \sigma b\) and \(f' b \xrightarrow{\text{fwd}} \sigma b\), we have \(f b = f' b\), implying \(f = f'\) since each \(f\) performs \(f b\) at most once. Thus we have \(w_i a \xrightarrow{\text{fwd}} f a\) and \(w'_i a \xrightarrow{\text{fwd}} f a\), which by Property (M.3) we have \(w_i a = w'_i a\), which again since we only have one write to \(\mathbb{A}[i]\) per abs write we have \(w_i = w'_i\).

**Property (F.5b).** If \(\sigma b_i\) terminated without \(\sigma r_i \xrightarrow{\text{fwd}} \sigma b_i\) then there exists \(w_i\) such that \(w_i \xrightarrow{\text{fwd}} \sigma\).

**Proof.** Since we do not have \(\sigma r_i \xrightarrow{\text{fwd}} \sigma b_i\), it is the case that \(\sigma b_i\) could not have observed any other \(\sigma' r_j\) since such \(\sigma'\) must have occurred before \(\sigma\) by virtual scanner total order Property (F.4a), implying \(\sigma' r_j \sqsubset \sigma r_j\), and \(\sigma r_j \triangleleft \sigma b_i\) by Property (F.4b), meaning we would contradict Property (M.2) if we had \(\sigma' r_j \xrightarrow{\text{fwd}} \sigma b_i\) by \(\sigma r_j\) occurring in between them. By Property (F.3b) this means that whatever \(\sigma b_i\) read, it cannot have been \(\perp\). By Property (F.3c), \(\sigma b_i\) must have observed some \(f b\), i.e., \(f b \xrightarrow{\text{fwd}} \sigma b_i\). This forwarding \(f\) must have read something with \(f a\) by the structure enforced by Property (F.4d), which by Property (F.3a) must be some \(w_i a\). Thus we have \(w_i a \xrightarrow{\text{fwd}} f a\) and \(f b \xrightarrow{\text{fwd}} \sigma b_i\), matching the instantiation of \(\xrightarrow{\text{fwd}}\).

**Property (F.5c).** If we have \(w_i \xrightarrow{\text{fwd}} \sigma\) then we have \(w_i a \triangleleft \sigma b\) without \(\sigma r_i \xrightarrow{\text{fwd}} \sigma b_i\).

**Proof.** Since \(w_i \xrightarrow{\text{fwd}} \sigma\) is instantiated as \(w_i a \xrightarrow{\text{fwd}} f a\) and \(f b \xrightarrow{\text{fwd}} \sigma b\), and since we have \(f a \sqsubset f b\) we have \(w_i a \xrightarrow{\text{fwd}} f a \sqsubset f b \xrightarrow{\text{fwd}} \sigma b\), implying \(w_i a \triangleleft \sigma b\). For proving absence of \(\sigma r_i \xrightarrow{\text{fwd}} \sigma b_i\), assume we have it, we derive a contradiction. By the instantiation of \(w_i \xrightarrow{\text{fwd}} \sigma\) we have \(f b \xrightarrow{\text{fwd}} \sigma b_i\), which by Property (M.3) gives us \(f b = \sigma r_i\), however that is impossible, since \(f b\) is a rep event of \(f\), while \(\sigma r_i\) is a rep event of \(\sigma\), thus a contradiction.

We have now satisfied Properties (F.4b) and (F.5c), so we satisfy the preconditions to use Lemma 3.1. Before we prove Properties (F.6) to (F.7b), we prove some helper lemmas. First, we prove that if a virtual scan \(\sigma\) observes a forwarding \(f\), then \(f\) must have observed \(\sigma\) on both in \(w_i x\) and \(f x\), meaning that the forward was concurrent to the scan.

**Lemma D.1.** If \(f b \xrightarrow{\text{fwd}} \sigma b\) then \(\sigma_{on} \xrightarrow{\text{fwd}} w_i x\) and \(\sigma_{on} \xrightarrow{\text{fwd}} f x\).

**Proof.** By \(f b \xrightarrow{\text{fwd}} \sigma b\), we know that \(f b\) successfully wrote to \(\sigma.B[i]\), thus we know by Property (M.7) that there exists some \(e_r\) that wrote to \(\sigma.B[i]\) such that \(e_r \xrightarrow{\text{fwd}} f b\) and \(e_r \xrightarrow{\text{fwd}} f b\). Also by the execution of \(f\) we know that both \(w_i x\) and \(f x\) observed some \(\sigma'_{on}\) of some \(\sigma'\), i.e., \(\sigma'_{on} \xrightarrow{\text{fwd}} w_i x\) and \(\sigma'_{on} \xrightarrow{\text{fwd}} f x\). To reach our goal, we prove by Property (F.4a) that we have \(\sigma = \sigma'\) by showing that we can neither have \(\sigma \sqsubseteq \sigma'\) nor \(\sigma' \sqsubseteq \sigma\).

- For \(\sigma \sqsubseteq \sigma'\), we have \(\sigma b_i \sqsubset \sigma' r_i\) by Property (RB.2) which leads to the cycle
  \[\sigma b_i \sqsubset \sigma' r_i \sqsubset \sigma'_{on} \xrightarrow{\text{fwd}} f x \sqsubset f b \xrightarrow{\text{fwd}} \sigma b_i\]
  which contradicts irreflexivity of \(\triangleleft\) (Lemma 2.1).

- For \(\sigma' \sqsubseteq \sigma\), we have \(\sigma' b_i \sqsubset \sigma r_i\) by Property (RB.2). By Property (M.4), we have \(\sigma r_i \triangleleft f b\) since if we instead had \(f b \triangleleft \sigma r_i\) we have \(\sigma r_i\) occurring in between \(f b \xrightarrow{\text{fwd}} \sigma b_i\), contradicting Property (M.2). Similarly by Property (M.4), we have \(e_r \triangleleft \sigma r_i\), since consider if we have \(\sigma r_i \triangleleft e_r\). By Property (RB.1) over \(f b \xrightarrow{\text{fwd}} f x\) and \(\sigma'_{off} \sqsubset \sigma' b_i\), we either have \(f b \triangleleft \sigma' b_i\), which lets us construct the following cycle contradicting irreflexivity of \(\triangleleft\) (Lemma 2.1)
  \[\sigma r_i \triangleleft e_r \xrightarrow{\text{fwd}} f b \sqsubset \sigma' b_i \triangleleft \sigma r_i\]
or we have $\sigma_{\text{off}} \subset f\tilde{x}$ which implies $\sigma_{\text{off}} \not< f\tilde{x}$ by $\sigma_{\text{off}} \not\leq \sigma_{\text{off}} \subset f\tilde{x}$, thus $\sigma_{\text{off}}$ occurs in between $\sigma_{\text{on}} \not\subset f\tilde{x}$, contradicting Property $(M^+2)$. Thus we have $e_r \not< \sigma_i$ and $\sigma_i \not< f\tilde{b}$, however this implies that $\sigma_i$ occurs in between $e_r \not\subset f\tilde{b}$, contradicting Property $(M^+2)$. □

**Lemma D.2.** If $w'_i \sqsubseteq \sqsubseteq \sigma$ then $\sigma_{\text{off}} \not< w'_i x$, and for any $f'$ executed by $w'_i$, we have $\sigma_{\text{off}} \not< f'$ and if $f'\tilde{b}$ was executed, then $f'\tilde{b} \not< \sigma\tilde{b}_i$.

**Proof.** For $w'_i \sqsubseteq \sqsubseteq \sigma$, there exists some $e$ and $w_j$ such that $w'_i \subset e \sqsubseteq w_j \not\subset \sigma$. We know all of these events are populated except for $\sigma$ since they occur to the left of some $\sqsubseteq$. By Lemma 3.1 and Property (RB.2), we have every rep event of $w'_i$ returns before some $e_r$ in $e$, and $e_r \not\subset w_ja$. We case split $w_j \not\subset \sigma$ into either $w_ja \not\subset \sigma a_j$, or $w_ja \not\subset \sigma b_j$.

- If $w_ja \not\subset \sigma a_j$ then every rep event of $w'_i$ happens before $\sigma_{\text{off}}$ since we have $e_r \not\subset w_ja \not\subset \sigma a_j \subset \sigma_{\text{off}}$ thus we have $\sigma_{\text{off}} \not< w_ja$ and $\sigma_{\text{off}} \not< f'x$ and $f'\tilde{b} \not< \sigma_{\text{off}} \not< \sigma b_i$ for all $f'$ executed by $w'_i$.

- If $w_ja \not\subset \sigma b_j$ then by Lemma D.1 we have $w_j \not\subset \sigma x$ and $w_j \not\subset \sigma f\tilde{x}$. By Property (F.3d) we have $\sigma_{\text{off}} \not< f\tilde{x}$, which in turn implies $\sigma_{\text{off}} \not< w_ja$ and $\sigma_{\text{off}} \not< f'\tilde{b}$, since every rep event in $w'_i$ occurs before $e_r$ and $e_r \not\subset w_ja \not\subset f\tilde{x} \not\subset f\tilde{b}$. By Property (RB.1) over $f'\tilde{b} \subset e_r$ and $\sigma_{\text{off}} \subset \sigma b_i$, we have $f'\tilde{b} \subset \sigma b_i$, since $\sigma_{\text{off}} \subset \sigma b_i$ contradicts $\sigma_{\text{off}} \not< f\tilde{x}$ by $\sigma_{\text{off}} \not< \sigma_{\text{off}} \not< e_r \not< w_ja \not< f\tilde{x}$, thus we satisfy all our goals.

□

**Lemma D.3.** If $\sigma_{\text{on}} \not\subset \sigma x$ with $f$ and $f'$ executed by $w_j$, where $f \subset f'$ and $\sigma_{\text{on}} \not\subset f\tilde{x}$ and $\sigma_{\text{on}} \not\subset f'\tilde{x}$ and $f'\tilde{b} \subset \sigma b_i$, then there exists some $f''$ such that $\sigma_{\text{on}} \not< f''\tilde{b}$ and $f''\tilde{b} \not< f\tilde{b}$ where $f''\tilde{b}$ successfully wrote to $\sigma.B[i]$.

**Proof.** Since we have $\sigma_{\text{on}} \not< w_ja$ and $f''\tilde{b} \not< \sigma b_i$, it is not possible for any $\sigma_{\text{off}}$ to occur in the intervals of $f$ and $f'$, therefore we can apply Lemma 4.3, thus by Property (F.3c) there exists some $f''$ such that $f''\tilde{b} \not< f\tilde{b}$ and $f''\tilde{b} \not< f'\tilde{b}$ where $f''\tilde{b}$ successfully wrote to $\sigma.B[i]$. By Property (RB.1) over $w_ja \subset f\tilde{b}$ and $f''\tilde{b} \subset f'\tilde{b}$, since $f''\tilde{b} \not< f\tilde{b}$ contradicts $f''\tilde{b} \subset f\tilde{b}$, we have $w_ja \subset f'\tilde{b}$, which lets us construct $\sigma_{\text{on}} \not< \sigma_{\text{on}} \not\subset w_ja \subset f''\tilde{b}$ implying $\sigma_{\text{on}} \not< f''\tilde{b}$, thus we satisfy all the goals.

□

**Property (F.6).** If $w'_i \sqsubseteq \sqsubseteq \sigma$ and $\sigma_{\text{on}} \not\subset \sigma b_i$ then $w_ja \not< \sigma a_i$.

**Proof.** By Lemma D.2 we have $\sigma_{\text{off}} \not< w'_i a$ and $\sigma_{\text{off}} \not< f'x$ and $\sigma_{\text{off}} \not< f'x$ and $f''\tilde{b} \not< \sigma b_i$ where $f'$ and $f''$ where executed by $w'_i$. Consider if we have $\sigma_{\text{on}} \not< w'_i x$ (and thus also $\sigma_{\text{on}} \not< f'x$ and $\sigma_{\text{on}} \not< f''x$), then by Property (F.3d) we have $\sigma_{\text{on}} \not\subset w'_i x$ and $\sigma_{\text{on}} \not\subset f'x$ and $\sigma_{\text{on}} \not\subset f''x$. By Lemma D.3 we know there exists some $f$ such that $\sigma_{\text{on}} \not< f\tilde{b}$ and $f\tilde{b} \not< f''\tilde{b}$ with $f\tilde{b}$ successfully writing to $\sigma.B[i]$. However, this means we have $\sigma_{\text{on}} \not< f\tilde{b} \not< f''\tilde{b} \not< \sigma b_i$, contradicting Property (M.2) by $f\tilde{b}$ occurring in between $\sigma_{\text{on}} \not\subset \sigma b_i$.

Thus we must have $\sigma_{\text{on}} \not< w'_i x$. By Property (RB.1) over $w'_i a \subset w'_i x$ and $\sigma_{\text{on}} \subset \sigma a_i$, since having $\sigma_{\text{on}} \not< w'_i x$ contradicts $\sigma_{\text{on}} \not< \sigma_{\text{on}} \subset w'_i x$, we must have $w_ja \not< \sigma a_i$, which is our goal, thus we are done.

□

**Properties (F.7a) and (F.7b).** If $w_j \not\subset \sigma a_i$ and either $w'_i \sqcap \sigma \not\subset \sigma a_i$ or $w'_i \not\subset w'_i a$ then $w'_i \not< w'_i a$.

**Proof.** Assume we have $w'_i \not< w'_i a$, we derive a contradiction. By Lemma 3.1 we have $w'_i \not< w'_i a$ from $w_i \not< w'_i a$, since both writes are populated by occurring to the left of some $\not\subset$. By the instantiation of $\not\subset$, there exists some $f$ such that $w_i \not\subset f\tilde{a}$ and $f\tilde{b} \not\subset \sigma b_i$. By Lemma D.1, we have $\sigma_{\text{on}} \not\subset w_i \not\subset w'_i$. We start with the $w'_i a \not< \sigma a_i$, case, where we can derive that $w'_i a \not< f\tilde{a}$:

$$w'_i a \not< \sigma a_i \subset \sigma_{\text{on}} \not\subset w'_i x \subset f\tilde{a}$$
which with \( w_i a < w_f a \), contradicts Property (M.2) by \( w_f a \) occurring in between \( w_i a \xrightarrow{rf} f a \).

For the \( w_f a \sqsubseteq w_f a \), \( \sigma \) case, we start by deriving \( f b \sqsubseteq w_f a \) by Property (RB.1) over \( f b \sqsubseteq f a \) and \( w_i a \sqsubseteq w_f a \), since we cannot have \( w_i a \sqsubseteq f a \) since that contradicts Property (M.2) by \( w_f a \) occurring in between \( w_i a \xrightarrow{rf} f a \). From this, we can derive that we have \( w_i a \sqsubseteq w_f a \) by \( w_i a \sqsubseteq f b \sqsubseteq w_f a \). By Lemma D.2 we have \( \sigma_{off} \notin w_f a \), which by Property (F.3d) means we have \( w_f a \sqsubseteq f b \), \( w_f a \), thus we know that \( w_f i_j \) and \( w_f i_2 \) will be executed, let us refer to these as \( f' \) and \( f'' \) respectively. Thus again by Lemma D.2 and Property (F.3d) we have \( \sigma_{on} \sqsubseteq w_f a \) and \( \sigma_{on} \sqsubseteq f' a \) and \( \sigma_{on} \sqsubseteq f'' a \).

By Lemma D.3 over \( f' \) and \( f'' \), there exists some \( f''' \) such that \( f''' a \notin f' b \) and \( f''' a \subseteq f'' b \). By Property (M.4), we either have \( f_b \sqsubseteq f''' b \) or \( f''' b \sqsubseteq f_b \). In the first case, we have

\[
f_b \sqsubseteq f''' b \sqsubseteq f'' b \preceq \sigma_b i
\]

which contradicts Property (M.2) by \( f''' b \) occurring in between \( f_b \sqsubseteq f'' b \). In the second case, if we have \( f''' b \sqsubseteq f_b \) then by Lemma 4.1, and \( f_b \) and \( f''' b \) being successful, we have \( f''' b \sqsubseteq f_b \), and this also follows in the case of \( f''' b = f_b \). By Property (RB.1) over \( w_f a \sqsubseteq f' b \) and \( f''' b \sqsubseteq f'' b \) we have \( w_f a \sqsubseteq f'' b \) since we cannot have \( f''' b \sqsubseteq f'' b \) by \( f''' b \notin f' b \). Thus we have

\[
w_i a \sqsubseteq w_i a \sqsubseteq f''' b \sqsubseteq f_b \sqsubseteq f_a
\]

which contradicts Property (M.2) by \( w_i a \) occurring in between \( w_i a \xrightarrow{rf} f a \). \( \square \)

**E ALGORITHMS 2 AND 3 SATISFY THE MULTI-WRITER FORWARDING SIGNATURE**

We will show the complete proofs that Algorithms 2 and 3 satisfy the multi-writer forwarding signature (Fig. 11). We start by repeating the proof that Algorithm 2 satisfies the signature.

**Lemma 4.6.** Every execution of Algorithm 2 satisfies the MWForwarding signature (Fig. 11).

**Proof.** Since this algorithm is single scanner, we can simply define the set of virtual scans to be the same as the set of abs scans and map each abs scan to itself, i.e., \( \Sigma = S \) and \( s^* = s \). Each rep event directly corresponds to their equivalent variant in Algorithm 2, except for \( \sigma_{on} \) and \( \sigma_{off} \) which are set as \( \sigma_{on} = \sigma_{on} = s_{on} \) and \( \sigma_{off} = \sigma_{off} = s_{off} \) respectively. Property (F.1) holds by each event being a subevent of itself (\( e \sqsubseteq e \)) and Property (F.4) holds since the algorithm is single-scanner. Properties (F.2) to (F.3c) and (F.4b) to (F.5b) holds directly by the structure of the algorithm and Property (F.3d) can be proven by Lemma C.1, the same lemma that established this property for Algorithm 1, since the relative structure of \( s_{on} \) and \( s_{off} \) is the same. \( \square \)

We next prove that Algorithm 3 satisfies signature. We first prove that the phases of logical scans can only happen in sequence.

**Lemma E.1.** Each \( v_{on} \) and \( v_{off} \) and \( v_{end} \) that successfully wrote to \( X \) are ordered such that there exists either \( v'_{on} \) or \( v'_{off} \) or \( v'_{end} \), where we have either \( v_{on} \xrightarrow{rf} v_{off} \xrightarrow{rf} v_{end} \) or \( v_{end} \xrightarrow{rf} v_{off} \xrightarrow{rf} v_{on} \), with the exception of the first \( v_{on} \), which observes the initial value of \( X \).

**Proof.** Looking at Algorithm 3, we see that the only operations writing to \( X \) are successful SC operations, which are only events \( v_{on} \) and \( v_{off} \) and \( v_{end} \). For either of these events to be successful, they need to observe the same write as the latest \( vx \) observed as per Property (M.7).

For \( v_{on} \) to be successful, we must have \( e_r \xrightarrow{rf} v_{on} \) and \( e_r \xrightarrow{rf} v_{on} \). For the algorithm to execute \( v_{on} \), the write \( e_r \) must have written 1 to x-phase, which can only be either the original initialization of \( X \) or \( v_{end} \). Similarly, for \( v_{off} \), it can only be executed if the write \( e_r \) wrote 2 to x-phase, thus \( e_r \) can only be \( v_{on} \), and for \( v_{end} \), it can only be executed if the write \( e_r \) wrote 3 to x-phase, thus \( e_r \) can only be \( v_{off} \). \( \square \)
We instantiate the write and forward rep events for the event structures of Fig. 10 with rep events of the same name. The structure of \texttt{write} and \texttt{forward} ensures that Properties (F.3a), (F.3c) and (F.4c) to (F.5b) are satisfied.

We instantiate virtual scans as a logical object corresponding to the rep events executed by \texttt{pushVS} constructing a complete logical scan. The interval of a virtual scan \(\sigma\) is instantiated as the smallest interval containing all its rep events. More formally, if we have \(v, v'\) and \(v''\) such that \(v_{on}(\rightarrow) v(\rightarrow) v'\) and \(v_{off}(\rightarrow) v''(\rightarrow) v_{ss}\) wrote to SS, then there exists a virtual scan \(\sigma\). Its rep events are instantiated using the rep events of \(v, v'\) and \(v''\) as follows and instantiate \(s^* = \sigma\) iff \(v_{ss}(\rightarrow) s_{ss}\).

\[
\begin{align*}
\sigma_1 & \triangleq vr_i \quad \sigma_{on} \triangleq v_{on} \quad \sigma_{on} \triangleq v(\rightarrow) \quad \sigma_i \triangleq va_i' \quad \sigma_{off} \triangleq v_{off}' \quad \sigma_{off} \triangleq v''(\rightarrow) \quad \sigma_{b} \triangleq vb_i''
\end{align*}
\]

We also let \(\sigma_{init} \triangleq >\) and \(\sigma_{ss} \triangleq v_{ss}'\) and \(\sigma_{end} \triangleq v_{end}'\) where \(v_{end}'\) was successful and \(v_{off}(\rightarrow) v_{ss}'\), however these do not need to be in the interval of \(\sigma\). We can view the events as forming the following structure.

\[
\begin{align*}
\sigma & \preceq v & v' & v'' \\
 vr_0 \cdots vr_{n-1} \subseteq v_{on}(\rightarrow) v(\rightarrow) va_0' \cdots va_{n-1}' \subseteq v_{off}(\rightarrow) v(\rightarrow) vb_0'' \cdots vb_{n-1}'' \subseteq v_{ss}'
\end{align*}
\]

Properties (F.2), (F.3b), (F.3d) and (F.4b) follows from the structure of virtual scans and the algorithm.

**Lemma E.2.** For every virtual scan \(\sigma\) such that \(\sigma_{end}\) have written to \(X\), there exists exactly one \(\sigma_{ss}\) such that \(\sigma_{off} \prec \sigma_{ss} \prec \sigma_{end}\).

**Proof.** By induction on the number of virtual scans performed. The base case corresponds to where \(\sigma\) is the very first virtual scan. We simultaneously consider the inductive step case, where the inductive hypothesis states that all virtual scans \(\sigma' \in \Sigma'\) has exactly one \(\sigma_{ss}'\) such that \(\sigma_{off}' \prec \sigma_{ss}' \prec \sigma_{end}'\), the inductive step follows that for \(\sigma\) proceeding the virtual scans of \(\Sigma'\), there exists exactly one \(\sigma_{ss}\) such that \(\sigma_{off} \prec \sigma_{ss} \prec \sigma_{end}\).

For the base case, since no other virtual scans has finished prior, we must have \(x.sync = ss.sync\) before \(\sigma_{off}\) has been executed by initialization of \(X\) and SS. For the inductive step, we have \(x.sync = ss.sync\) before \(\sigma_{off}\) executes, since it follows from each \(\sigma' \in \Sigma'\) having \(\sigma_{ss}' \prec \sigma_{end}'\) and there only existing one \(\sigma_{ss}'\) per \(\sigma'\), since after \(\sigma_{end}'\) has executed, \(x.sync = ss.sync\) by \(\sigma_{ss}'\) negating \(ss.sync\) and \(\sigma_{end}'\) negating \(x.sync\), both exactly once.

With \(\sigma_{end}\) successfully writing, we know that the code of phase 3 (lines 31-43 of Algorithm 3) must have been executed. Let \(\sigma_{end} = v_{end}\), we also know that \(v_{ss}\) returned \(true\), since otherwise \(\sigma_{end}\) must have failed to write. If the value of \(ss.sync\) read by \(v_{ss}\) is different from the value of \(x.sync\) read by \(\sigma_{off}\), then there must exist some \(v_{ss}'\) that negated \(ss.sync\) before \(v_{ss}\) (and thus before \(\sigma_{end}\)), but after \(\sigma_{off}\), which fits our goal.

Otherwise, if \(ss.sync = x.sync\), then \(v_{ss}\) will be executed. If \(v_{ss}\) successfully wrote, then we trivially satisfy our goal. Otherwise, by Lemma 4.2, there has to exist some \(v_{ss}'\) such that \(v_{ss}' \prec v_{ss}\) and \(v_{ss} < v_{ss}\). Since we must have \(v_{off}(\rightarrow) v(\rightarrow) v_{ss}''(\rightarrow) v_{ss}'\) for some \(v_{ss}''\), consider if we have \(\sigma_{off} = v_{ss}''\), then we have \(\sigma_{off} < v_{ss}' < \sigma_{end}\), satisfying our goal. Consider instead we have \(\sigma_{off} = v_{ss}'\), by Property (M.4), we either have \(\sigma_{off} < v_{ss}''\) or \(v_{ss}'' < \sigma_{off}\), where if we have the former then we must have \(\sigma_{off} < \sigma_{end} < v_{ss}'' < v_{ss}\) by Lemma E.1, inducing the cycle \(v_{ss} \subseteq \sigma_{end} < v_{ss}'' < v_{ss}\) contradicting irreflexivity of \(<\), thus we must have the latter. By \(v_{ss}'' < \sigma_{off}\), it must be the case that \(v_{ss}'\) correspond to some earlier executed virtual scan, which in the base case contradicts \(\sigma\) being the first virtual scan, while in the inductive step \(v_{ss}'\) corresponds to some \(\sigma' \in \Sigma'\), it must be the
case that \( v'_{ss} < \sigma'_{end} \ll \sigma_{off} < v_{ss} \) by the inductive hypothesis and Lemma E.1, which contradicts \( v'_{ss} \not\prec v_{ss} \).

To efficiently reason about \textsc{pushvs}, we define the intervals of the phases. For an execution \( v \) of \textsc{pushvs}, \( v.p_1 \) and \( v.p_2 \), and \( v.p_3 \) are the intervals of phases 1, 2, and 3. Each phase interval \( p \) of \( v \) has a start event \( p.se \) which is some \( v_x \) and an end event \( p.\text{ee} \) which is a SC event of \( X \), where start and end events need to satisfy \( p.se \ll p.\text{ee} \). For \( p_1 \) we have \( p_1.\text{ee} = v_{on} \) with either \( p_1.se \) observing the initial value of \( X \) or \( \sigma_{end} \ll p_1.\text{ee} \); for \( p_2 \) we have \( p_2.\text{ee} = v_{off} \) with \( \sigma_{on} \ll p_2.\text{ee} \); and for \( p_3 \) we have \( p_3.\text{ee} = v_{end} \) with \( \sigma_{off} \ll p_3.\text{ee} \). Each phase \( p \) has an advancement event, which we encode with \( p.ae \), which is the rep event successfully writing to \( X \) which advances the logical scanner to the next phase. For some \( \sigma \), we have \( p_1.ae = \sigma_{on} \) and \( p_2.ae = \sigma_{off} \) and \( p_3.ae = \sigma_{end} \).

**Lemma E.3.** For each phase \( p \), there exists a \( p.ae \) such that \( p.ae \not\prec p.se \) and \( p.ae \ll p.\text{ee} \).

**Proof.** By Lemma 4.2, there exists some successful writing event \( e_r \) of \( X \) such that \( e_r \not\prec p.se \) and \( e_r \ll p.\text{ee} \). If \( e_r \not\prec p.ae \), we are done, otherwise by Lemma E.1 we know that \( \sigma_{on} \) and \( \sigma_{off} \) and \( \sigma_{end} \) occur in order, so for example in the case of \( p = p_2 \), we have \( \sigma_{on} \ll p_2.\text{ee} \), by Property (M*4) we must have \( \sigma_{on} \ll e_r \), since \( e_r \ll \sigma_{on} \) contradicts \( e_r \not\prec p.se \). Thus by Lemma E.1, there exists some \( \sigma_{off} \ll p_2.ae \) in between \( \sigma_{on} \) and \( e_r \), which satisfies \( \sigma_{off} \ll p.\text{ee} \) by \( \sigma_{off} \ll e_r \ll p.\text{ee} \) and \( \sigma_{off} \not\prec p.se \) since \( \sigma_{off} \ll p.se \) contradicts Property (M*2) by \( \sigma_{off} \) occurring in between \( \sigma_{on} \ll p_2.\text{ee} \). The other phases follow by similar arguments.

**Lemma E.4.** For any terminated \textsc{pushvs} execution \( v \), there exists some \( \sigma \) such that for \( v_x \) being the first rep event of \( v \) and \( e_r \) being the last, then and \( \sigma_{end} \not\prec v_x \) and \( \sigma_{end} \ll e_r \).

**Proof.** Consider whether \( p_3 \) was executed or not by \( v \). If \( p_3 \) was executed, then by Lemma E.3 we satisfy our goal. Otherwise, if it was not executed, this means either \( p_1 \) or \( p_2 \) were executed, i.e., there is some \( \sigma_{end} \) or \( \sigma_{on} \) setting \( x.phase \), thus there is another chance for \( p_3 \) to run. If again it was not run, we are either in \( p_1 \) or \( p_2 \) by \( \sigma_{end}/\sigma_{on} \), by Lemma E.3 the execution of virtual scans must have progressed, and by Lemma E.1 there must be some \( \sigma''_{end} \) between \( \sigma_{end}/\sigma_{on} \) and \( \sigma'_{end}/\sigma'_{on} \), where \( \sigma''_{end} \) may be equal to \( \sigma'_{end} \). Thus \( \sigma''_{end} \) satisfies our goal.

**Lemma E.5.** For any terminated \textsc{pushvs} executions \( v \) and \( v' \) where \( v \sqsubseteq v' \), there exists some \( \sigma \) and \( \sigma' \) such that for \( v_x \) being the first rep event of \( v \), then \( \sigma_{end} \not\prec v_x \) and \( \sigma_{end} \ll \sigma'_{end} \).

**Proof.** By Lemma E.4 over \( v \) and \( v' \), we have \( \sigma_{end} \not\prec v_x \) and \( \sigma_{end} \ll e_r \) and \( \sigma'_{end} \not\prec v'_x \) and \( \sigma'_{end} \ll e'_r \) where \( e_r \) and \( e'_r \) are the last events of \( v \) and \( v' \) respectively. By Property (M*4), we must have \( \sigma_{end} \ll \sigma'_{end} \), since if we instead had \( \sigma'_{end} \not\prec \sigma_{end} \), we contradict \( \sigma'_{end} \not\prec v'_x \) by having \( \sigma'_{end} \ll \sigma_{end} \ll e_r \ll v'_x \).

**Property (F.1).** For any abs scan \( s \), we have \( s^* \sqsubseteq s \).

**Proof.** With Lemma E.5, we can show for \( s \) that there exists some \( \sigma'_{end} \ll \sigma''_{end} \) in between \( sv_1 \) and \( sv_2 \) such that \( \sigma'_{end} \not\prec v_x \) where \( v_x \) is the first rep event of \( sv_1 \). For \( \sigma = s^* \), we have \( \sigma'_{end} \ll \sigma_{init} \) since some new logical scan must have started after \( \sigma'_{end} \) by \( \sigma''_{end} \) occurring after it, and by the instantiation of \( [-]^{n} \) we have \( \sigma_{init} \ll \sigma_{init} \ll s_{ss} \), thus for all \( i \) we have \( \sigma'_{end} \ll \sigma_{init} \ll \sigma_{init} \ll s_{bh} \ll s_{ae} \ll s_{ss} \). Our goal is to establish \( s.\text{start} \leq \sigma_{init}.\text{start} \) and \( s_{bh}.\text{end} \leq s.\text{end} \). We show this by showing that we neither have \( s.\text{start} > \sigma_{init}.\text{start} \) nor \( s_{bh}.\text{end} > s.\text{end} \).

- Assume we have \( s.\text{start} > \sigma_{init}.\text{start} \). We derive a contradiction. By definition of \( \sigma_{init} \not\prec \sigma_{init} \), we have \( \sigma_{init}.\text{end} < \sigma_{init}.\text{start} \), which implies \( \sigma_{init}.\text{end} < s.\text{start} \) and thus \( s_{bh} \not\prec s \). This in turn implies \( s_{bh} \not\prec v_x \) since \( v_x \not\prec s \) which contradicts \( \sigma'_{end} \not\prec v_x \) by \( \sigma'_{end} \not\prec \sigma_{init} \not\prec v_x \).
• Assume we have $\sigma b_j.\text{end} > s.\text{end}$, we derive a contradiction. By definition of $\sigma b_j \sqsubseteq \sigma_{ss}$, we have $\sigma b_j.\text{end} < \sigma_{ss}.\text{start}$, which implies $s.\text{end} < \sigma_{ss}.\text{start}$ and thus $s \sqsubset \sigma_{ss}$. This in turn implies $s_{ss} \sqsupset \sigma_{ss}$ since $s_{ss} \subseteq s$, which contradict irreflexivity of $\prec$ by $s_{ss} \sqsupset \sigma_{ss} \Rightarrow s_{ss}$.

**Property (F.3d).** For any $e_r$ reading $X$, we have $s_{on} \nleq_{rf} e_r$ iff we have $s_{on} \prec e_r$ and $s_{off} \not\prec e_r$.

**Proof.**

$(\Rightarrow )$ $s_{on} \nleq_{rf} \sigma \beta \Rightarrow$ trivially implies $s_{on} \prec \sigma \beta$, and if we were to have $s_{off} \prec \sigma \beta$ we contradict Property (M.2) since $s_{off}$ would happen in-between $s_{on}$ and $e_r$.

$(\Leftarrow )$ By the fact that only virtual scans change $X$ and by Lemma E.1 forcing these changes to be of an order of form $s_{on} \nleq_{rf} \sigma \beta \nleq_{rf} \sigma_{end} \nleq_{rf} s_{on} \ldots$, we know that any write to $X$ other than $s_{on}$ and $s_{off}$ either happens before $s_{on}$ or after $s_{off}$. Since $\sigma \beta \beta$ has to observe something when it finishes, it must observe $s_{on}$, since any other observation leads to a contradiction.

**Property (F.4a).** For two distinct virtual scans $\sigma$ and $\sigma'$, we either have $\sigma \sqsubseteq \sigma'$ or $\sigma' \sqsubseteq \sigma$.

**Proof.** By the instantiation of $\sigma$ and Lemma E.1, we have that each virtual scan must be separated by some $\sigma_{ss}$, i.e., we either have $\sigma \prec \sigma' \text{init}$ or $\sigma \prec \sigma \text{init}$. By Lemma E.2, we know that for each $\sigma$ there is exactly one $\sigma_{ss}$ occurring before $\sigma_{end}$, thus we have either $\sigma_{ss} \prec \sigma_{end} \prec \sigma_{\text{init}}$ or $\sigma_{ss} \prec \sigma_{\text{end}} \prec \sigma_{\text{init}}$. If we have $\sigma_{ss} \prec \sigma_{end} \prec \sigma_{\text{init}}$ then we have $\sigma b_i \sqsubset \sigma_{ss} \prec \sigma_{\text{end}} \prec \sigma_{\text{init}} \sqsubset \sigma'_{\text{ss}}$, and dually for $\sigma'_{\text{ss}} \prec \sigma'_{\text{end}} \prec \sigma_{\text{init}}$, thus we either have $\sigma b_i \sqsubset \sigma'_{\text{ss}}$ or $\sigma b_i \sqsubset \sigma'_{\text{s}}$ for each $i$, which matches the defined bounds of $\sigma$, i.e., we either have $\sigma \sqsubset \sigma'$ or $\sigma' \sqsubset \sigma$, which is our goal.

### F. ALGORITHM 4 SATISFIES THE SNAPSHOT SIGNATURE

We show that Algorithm 4 satisfies the Snapshot signatures (Fig. 5). We employ virtual scans as a simplification mechanism that allows us to elide write views from consideration when establishing Snapshot. A virtual scan will denote a scan that detected no change in the array, i.e., it returned at line 21. Formally, a virtual scan $\sigma$ consists of rep events $\sigma a_i$ and $\sigma b_i$ for each $i$, where both $\sigma a_i$ and $\sigma b_i$ observe the same write (i.e., no change detected), and every $\sigma a_i$ occurs before any $\sigma b_i$:

$$
\forall \sigma, i. \quad \exists w_i. w_i a \nleq_{rf} \sigma a_i \land w_i a \nleq_{rf} \sigma b_i \quad (A.1 \text{ revisited})
$$

$$
\forall \sigma, i, j. \quad \sigma a_i \sqsubseteq \sigma b_j \quad (A.2 \text{ revisited})
$$

To each abs scan, we can associate a virtual scan as follows. If the abs scan terminated because it detected no change, then it immediately is a virtual scan. If the abs scan terminated by returning a view from a write, i.e., returned at line 16, then that view itself is an abs scan which can, recursively, be associated with a virtual scan. We formally define this by $[-]^{*}$ as follows:

$$
\begin{align*}
\sigma^{*} & \triangleq \begin{cases} 
\sigma & \text{where} \quad \exists k. \quad \forall i. \quad \sigma a_i = s_k a_i \land \sigma b_i = s_k b_i \\
\sigma \beta \beta & \text{where} \quad \exists w_{i}' , w_{i}'' , j, k. \quad j < k \land w_{i}'' \nleq_{rf} w_{i}' \leq_{rf} w_i \land \\
& \quad w_{i}'' a \nleq_{rf} s_j a_i \land w_{i}'' a \nleq_{rf} s_k a_i \land w_i a \nleq_{rf} s_k b_i
\end{cases}
\end{align*}
$$

We instantiate $\sigma^{*}$ to be the set of all writes $w_i \in W_i$ where $w_i a$ is defined; these are the writes that executed their effect. We also instantiate the visibility relations $\prec$ and $\nleq_{rf}$ as follows, using the helper relation $\nleq_{rf}$.

$$
\begin{align*}
w_i \nleq_{rf} \sigma & \triangleq \quad w_i a \nleq_{rf} \sigma a_i \\
w_i \nleq_{rf} s & \triangleq \quad w_i \nleq_{rf} s^{*} & < \triangleq \quad \nleq_{rf}
\end{align*}
$$

In English, the scan $s$ reads from, and also observes, $w_i$ iff there is an appropriate rep event in the virtual scan $s^{*}$ that reads from $w_i a$ at the level of rep events.

**Lemma F.1.** $[-]^{*}$ is a well-founded recursive definition.
Proof. The first case of $[-]^*$ is immediately well-founded. For the second case it is well-founded because the ending time of the scan $w_{s'}$ on the right is smaller than the ending time of $s$ on the left, and the ending times are bounded from below by 0. To see that the ending time of $w_{s'}$ is below that of $s$, suppose otherwise. Then it must be $s \subseteq w_{a}$, because from the code of write we have that $w_{s'} \subseteq w_{a}$. In particular, $s_{k}b_{i} \subseteq w_{a}$ for every $k$. But, because $w_{a}$ is observed by some rep read in $s$, we also have $w_{a} \not\subseteq s_{k}b_{i}$ for some $k$. Thus, we have an event $s_{k}b_{i}$ that terminated before the event $w_{a}$ that it observes, which contradicts Property (V.1).

\[\square\]

Lemma F.2. For each scan $s$, we have $s^* \subseteq s$.

Proof. If $s^* = \sigma$ corresponds to the first case, then this is trivial, since $\sigma$ consists of the selected rep events of $s$, i.e., $\sigma_{a} = s_{k}a_{i}$ and $\sigma_{b} = s_{k}b_{i}$. Otherwise, $s^* = (w_{s'})^*$ corresponds to the second case for some write view $w_{s'}$. By recursion on the definition of $[-]^*$ it must be $s^* = (w_{s'})^* \subseteq w_{s'}$, so it suffices to show $w_{s'} \subseteq s$. By definition, if we have $s.start \leq w_{s}.start$ and $w_{s}.end \leq s.end$ then we have $w_{s} \subseteq s$. During the scan, at least three different writes must have been observed: $w_{a} \not\subseteq s_{k}b_{i}$ and $w_{a} \not\subseteq s_{j}a_{i}$ where we have $w_{s''} \subseteq w_{s'} \subseteq w_{a}$ by the algorithm being single-writer. If we have $w_{a} \subseteq s_{j}a_{i}$ we contradict Property (M.2) for $w_{s''}a \not\subseteq s_{j}a_{i}$ since we have $w_{s''} \subseteq w_{a} \subseteq s_{j}a_{i}$, thus we have $w_{a} \not\subseteq s_{j}a_{i}$, which by definition gives us $s_{j}a_{i}.start \leq w_{a}.end$. Similarly, if we have $w_{a} \subseteq s_{j}a_{i}$ we contradict Property (M.2) for $w_{s''}a \not\subseteq s_{j}a_{i}$ since we have $w_{s''} \subseteq w_{a} \subseteq s_{j}a_{i}$, thus we have $s_{k}b_{i} \not\subseteq w_{a}$, which by definition gives us $w_{a}.start \leq s_{k}b_{i}.end$. Together, we have

\[s.start < s_{j}a_{i}.start \leq w_{a}.end < w_{s}.start\]

and

\[w_{s}.end < w_{a}.start \leq s_{k}b_{i}.end < s.end,\]

which gives us $s.start \leq w_{s}.start$ and $w_{s}.end \leq s.end$, i.e., $w_{s} \subseteq s$.

Proof (V.1). If $e < e'$ then we cannot have $e' \subseteq e$.

Proof. If $e < e'$ and $e' \subseteq e$ then $e < e'$ can only hold if $e \not\subseteq e'$, i.e., we have $e = w_{i}$ and $e' = s$ and $\sigma = s^*$ such that $w_{i}a \not\subseteq \sigma_{a}$. We have that $e' \subseteq e$ can only be $e' \subseteq e$, since $e$ and $e'$ cannot be equal events by $e$ being a write and $e'$ being a scan. But then, by Property (RB.2) and Lemma F.2, from $e' \subseteq e$ we derive $\sigma_{a} \subseteq w_{a}$, which contradicts Property (V.1) for registers.

Property (S.1). For every terminated scan $s$ and index $i$, there exists some write $w_{i}$ such that $w_{i} \not\subseteq s$ and $w_{i}.in = s.out[i]$.

Property (S.3). If $w_{i} \not\subseteq s$ and $w_{i}' \not\subseteq s$ then we must have $w_{i} = w_{i}'$.

Proof. If $w_{i} \not\subseteq s$ and $w_{i}' \not\subseteq s$, then $w_{a}$ and $w_{a}'$ are both observed by some $\sigma_{a}$, i.e., the same read. Thus, by Property (M.3), $w_{a} = w_{a}'$, and since each write executes $w_{a}$ only once by Property (ES.2), it must be $w_{i} = w_{i}'$.

Proof. The structure of the algorithm ensures that at every return of scan, writes of every memory cell are observed and returned.

Property (S.4). For two distinct writes $w_{i}$ and $w_{i}'$, we have either $w_{i} \not\subseteq w_{i}'$ or $w_{i}' \not\subseteq w_{i}$.

Proof. Holds by the assumption that the algorithm is single-writer.

Property (S.5). Every terminated write is effectful, i.e., $T(W_{i}) \subseteq W_{i}$.

Proof. Since each terminated write must have $w_{a}$ defined, which is how we define the set $W_{i}$, thus every terminated write must be in some $W_{i}$.
Lemma F.3. If $w_i \prec s$ then $w_i \prec s^*$.

Proof. This proof follows similarly to the proof of Lemma B.1. By viewing $w_i \prec s$ as a chain of $\prec_1$, we have $w_i \prec_1 \ldots \prec_1 s$. By induction on the length of the chain. In our base case we have $w_i \prec_1 s$, which splits into cases $w_i \sqsubseteq s$ and $w_i \prec s$, where the latter can only be $w_i \xrightarrow{rf} s$. For $w_i \sqsubseteq s$, by Property (RB.2) and Lemma F.2, we have $w_i \sqsubseteq s^*$. For $w_i \xrightarrow{rf} s$, we have $w_i \xrightarrow{rf} s^*$ directly from the definition of $\xrightarrow{rf}$. The inductive step is similar. □

Lemma F.4. Choose some $j$. For events $e, e' \in \bigcup_j W_i \cup \Sigma$ where $e'$ is either a write $e' = w'_i$ such that $w'_ia$ is defined or a virtual scan $e' = \sigma'$, if $e \prec e'$ holds then there exists $e_r$ and $e'_r$ such that $e_r \prec e'_r$, where if $e = w_i \in W_i$ then $e_r = w_i a$ and if $e = \sigma \in \Sigma$ then $e_r = \sigma b_j$, and dually for $e'$ and $e'_r$.

Proof. This proof follows similarly to the proof of Lemma 3.1. We can view $e \prec e'$ as a chain of $\prec_1$ relations connecting $e$ and $e'$; that is, $e \prec_1 \ldots \prec_1 e'$, where by definition $\prec_1 = (\sqsubseteq \cup \prec) = (\sqsubseteq \cup \xrightarrow{rf})$. It is easy to see that each abs event in this chain is populated with a rep event. Indeed, if an abs event’s relation $\prec_1$ to its successor is realized by $\sqsubseteq$, then the event is terminated and all of its rep events are executed (and by the structure of the algorithm, each abs event must have at least one rep event). Alternatively, if the abs event’s relation $\prec_1$ to its successor is realized by $\prec$, then the event must be populated as per the definitions of reads-from. This leaves out $e'$, which has no successor, but $e'$ is populated by either being a write, thus having $w_i a$ defined by assumption, or being a virtual scan, which are first defined when a scan returns.

Now the proof is by induction on the length of the above chain. The base case is when $e \prec_1 e'$. If $e \sqsubseteq e'$, the proof is by Property (RB.2) giving us $e_r \sqsubseteq e'_r$ where if $e = w_i \in W_i$ then $e_r = w_i a$ and if $e = \sigma \in \Sigma$ then $e_r = \sigma b_j$, and dually for $e'$ and $e'_r$. For $e \xrightarrow{rf} e'$, we have $e = w_i$ and $\sigma' = e'$ with $w_i a \xrightarrow{rf} \sigma' a_i$. By Property (A.2) we construct $w_i a \xrightarrow{rf} \sigma' a_i \sqsubseteq \sigma' b_j$ which gives us our goal $w_i a \prec \sigma' b_j$. The inductive step is similar, with the addition that we need transitivity of $\prec$ to join events. □

Property (S.2). If $w_i \xrightarrow{rf} s$ then there does not exist a write $w'_i$ such that $w_i \prec w'_i \prec s$ holds.

Proof. Given $w_i \xrightarrow{rf} s$, we assume there exists $w'_i$ such that $w_i \prec w'_i \prec s$, we derive a contradiction. By assumption, we have $w_i \xrightarrow{rf} \sigma$ for $\sigma = s^*$. By Lemma F.3 and $w_i \prec w'_i \prec s$, we can derive $w_i \prec w'_i \prec \sigma$, and then by Lemma F.4, also $w_i a \prec w'_i a \prec \sigma b_i$. By the definition of $\xrightarrow{rf}$ and Property (A.1), $w_i \xrightarrow{rf} \sigma$ implies $w_i a \xrightarrow{rf} \sigma a_i \sqsubseteq \sigma b_i$. But then $w'_i a$ occurs between $w_i a$ and $\sigma b_i$, contradicting Property (M.2) and $w_i a \xrightarrow{rf} \sigma b_i$. □

Property (S.6). If $w_i, w_j \xrightarrow{rf} s$ and $w'_i, w'_j \xrightarrow{rf} s'$ with $w_i \prec w'_i$, then we cannot have $w'_j \prec w_j$.

Proof. Given $w_i, w_j \xrightarrow{rf} s$ and $w'_i, w'_j \xrightarrow{rf} s'$, with $w_i \prec w'_i$, we show that also having $w'_j \prec w_j$ lead to a contradiction. We unfold the assumptions $w_i, w_j \xrightarrow{rf} s$ and $w'_i, w'_j \xrightarrow{rf} s'$ into $w_i, w_j \xrightarrow{rf} \sigma$ and $w'_i, w'_j \xrightarrow{rf} \sigma'$, where $\sigma = s^*$ and $\sigma' = s'^*$. By the definition of $\xrightarrow{rf}$ and Property (A.1), from $w_j \xrightarrow{rf} \sigma$, we have $w_j a \xrightarrow{rf} \sigma a_j, \sigma b_j$, and similarly for $w'_j \xrightarrow{rf} \sigma'$. By Property (A.2) we have $\sigma a_j \sqsubseteq \sigma b_j$ and $\sigma' a_i \sqsubseteq \sigma' b_j$. By Property (RB.1), it is either $\sigma a_j \sqsubseteq \sigma' b_j$ or $\sigma' a_i \sqsubseteq \sigma b_j$. In the first case (the second is symmetric) we can construct $w_j a \prec w_j a \prec \sigma' b_j$ contradicting Property (M.2) by $w_j a$ occurring between $w'_j a \xrightarrow{rf} \sigma' b_j$. □