Data Hiding Using Sequential Hamming + k with m Overlapped Pixels

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Abstract

Recently, Kim et al. introduced the Hamming + k with m overlapped pixels data hiding (Hk_mDH) based on matrix encoding. The embedding rate (ER) of this method is 0.54, which is better than Hamming code HC (n, n − k) and HC (n, n − k) +1 DH (H1DH), but not enough. Hamming code data hiding (HDH) is using a covering function COV(1, n = 2k − 1, k) and H1DH has a better embedding efficiency, when compared with HDH. The demerit of this method is that they do not exploit their space of pixels enough to increase ER. In this paper, we increase ER using sequential Hk_mDH (SHk_mDH ) through fully exploiting every pixel in a cover image. In SHk_mDH, a collision maybe happens when the position of two pixels within overlapped two blocks is the same. To solve the collision problem, in this paper, we have devised that the number of modification does not exceed 2 bits even if a collision occurs by using OPAP and LSB. Theoretical estimations of the average mean square error (AMSE) for these schemes demonstrate the advantage of our SHk_mDH scheme. Experimental results show that the proposed method is superior to previous schemes.

Keywords: Data Hiding, Matrix Encoding, Hamming DH (HDH), Hamming+1 DH (H1DH), Hk_mDH
1. Introduction

Secure communication is often accomplished through cryptography. However, for secret communication, steganography is more suitable than cryptography. The reason is the purpose of steganography is that the existence of hidden data is not detected by the adversary [1]. This technique uses a communication channel to transmit secretly after concealing messages on various media and it is called data hiding (DH) [2]. DH is to embed secret bits by slightly modifying the cover media. In fact, images contain a lot of redundant bits, making them very suitable for delivering secret messages. Meanwhile, techniques for detecting hidden data in images, i.e., steganalysis [3, 4], have also been actively researched. Therefore, most researches on DH have been dedicated on minimizing the damage of an image while allowing sufficient messages to be concealed in images.

The method based on Least Significant Bit (LSB) is easily implemented by increasing or decreasing the values of pixels by only one to match the pixels of the image with the bits of the data. Meanwhile, in the aspect of image quality based on LSB, the optimal pixel adjustment process (OPAP) [5] is a great help for improving peak signal to noise ratio (PSNR) of stego image. Using only LSB, it is impossible to guarantee the quality of a stego image, because it could provoke distortion. Recently, Yang et al. [6] adopted LSB and OPAP with a cover image generated neighbor mean interpolation to improve visual quality in DH.

Matrix encoding is a method to embed or extract the message by using syndrome. Crandall [7] proposed a DH called matrix encoding, Westfeld implemented the F5 algorithm [8] using Hamming code, and Bierbrauer et al. [9] analyzed the possibility of this method. Willems [10] introduced a binary embedding method and a ternary embedding method using Hamming code and Golay code respectively. The F5 algorithm is a method that conceals $k$ bits into $(2^k - 1)$ pixels by using modifying the only one bit. F5 has higher embedding efficiency (EE) (which is the average number of embedded bits per change) [11] than the simple LSB method. Since then, many kinds of research based on matrix encoding have been developed inspired by the F5 [12].

Zhang [13] and Fridrich [14] proposed ternary Hamming code that each secret digit in a $(2n+1)$-ary notational system is carried by $n$ cover pixels by increasing or decreasing only one bit. Mielikainen [15] devised LSB matching revised method to achieve the same payload by modifying fewer pixels to the cover image than a simple LSB-based method. In Mielikainen’s scheme, the embedding procedure performs for two cover pixels, $x_i$ and $x_{i+1}$, at a time. The value of the $i^{th}$ message bit $m_i$ is equal to the LSB of the $i^{th}$ pixel $y_i$. The value of the $i + 1^{th}$ message bit $m_{i+1}$ is a function (LSB ($\lfloor y_i/2\rfloor + y_{i+1}$)) of $y_i$ and $y_{i+1}$. Zhang et al. [11] proposed Hamming + 1 DH (H1DH) embedding one more bit than COV (1, $2^k - 1$, $k$) with using one additional pixel. In H1DH, the performance improves by employing the OPAP. Recently, Kim and Yang [16] proposed a method of overlapping three pixels to improve ER with COV (1, $2^k - 1$, $k$), and the ER was increased to about 0.12 compared to HDH.

In [16], the second covering function COV (1, $2^k - 1$, $k$) uses a second LSB. Therefore, it causes a high distortion. Hamming + k (HkDH) [17] is a method to embed $2k$ in a fully overlapped virtual block (LSB $\oplus$ 2LSB) by using COV (1, $2^k - 1$, $k$), and OPAP and LSB were used for optimization. Therefore, HkDH has achieved high ER. Afterward, Kim et al. [18] extended HkDH to $m$ overlapped pixels data hiding (Hk_mDH) to embed $2k$ bits in $(2^{k+1} - m - 2)$ pixels by modifying at most 3 bits. But, if the values of two COV (1, $2^k - 1$, $k$) are the same, a collision problem may occur, which may inevitably cause a mean square error. Finally, the Hk_mDH achieves better embedding rate compared with
HkDH. The problem of HkDH is that the number of pixels that need to be excavated to embed many bits still remains. In this paper, we introduce sequential HkDH (SHkDH) to solve the matter of HkDH. SHkDH has a way of improving the performance of the ER without significantly degrading the quality of the cover image.

The rest of this paper is organized as follows. Section 2 concisely reviews Hamming code, HDH, and Kim and Yang’s DH, and HkDH. In Section 3, we present SHkDH. In addition, the AMSE of SHkDH is theoretically derived. In Section 4, the proposed SHkDH is tested using some images. A comparison with the HDH, H1DH, Kim and Yang’s DH, and HkDH is also provided. In addition, we apply RS steganlysis to SHkDH. Finally, this paper concludes in Section 5.

2. Related Work

2.1 Hamming Code

HC (n, k) codes are a single error correction linear code with dmin = 3 and parity check matrix H of k × n (= 2^k − 1), where k is the number of information bits and (n − k) is the number of parity bits. The parity check matrix H for k = 3 is

\[
H = \begin{bmatrix}
0 & 0 & 0 & 1 & 1 & 1 & 1 \\
0 & 1 & 1 & 0 & 0 & 1 & 1 \\
1 & 0 & 1 & 0 & 1 & 0 & 1 \\
\end{bmatrix}
\]

The n-bit column vectors \(x_b\) and \(y_b\) are LSB of pixels \(x = (x_1, \ldots, x_{n-k})\) and \(y = (y_1, \ldots, y_n)\), where \(x_b\) and \(y_b\) will be denoted \(F_2^{n-k}\) and \(F_2^n\), respectively.

Let \(y_b\) is obtained from an information bits \(x_b\) via \((n-k) \times n\) generator matrix \(G\), i.e., \(y_b = x_b \cdot G\). For any \(y_b\), the vector \(S = Hy_b^T(\in F_2^k)\) is called the syndrome of \(y_b\). For each syndrome \(S\), the set \(\hat{C}(S) = \{y_b \in F_2^n \mid Hy_b = S\}\) is called a coset. Note that \(\hat{C}(0) = \hat{C}\). For any syndrome \(S\), let \(\text{dec}(S)\) be the integer. For any non-zero syndrome \(S\), the vector \(e(S) = (0, \ldots , 1, \ldots, 0)\) with 1 at the \(\text{dec}(S)\)-th place, i.e., \(He(S) = S\). To embed \(\delta (\in F_2^k)\) message bits, replace \(y\) with \(\hat{y} = y_b \oplus e(\delta - Hy_b)\). The receiver extracts the message from \(\hat{y}\) by Eq. (2).

\[
\begin{align*}
\{\hat{y} \cdot H^T &= (y_b \oplus e) \cdot H^T = y_b \cdot H^T + \delta - y_b \cdot H^T \\\n\delta &= \delta
\}
\end{align*}
\]

In order to embed \(\delta\) message bits in each subset by making at most one embedding change, we divide the cover image into \(N/n(= 2^k − 1)\) subsets, each consisting of \(n\) pixels. Suppose that there is one error bit occurred in \(\hat{y}\) (say the 6-th bit from left), i.e., \(e = (e_1, \ldots, e_7) = (0000010)\). The syndrome \(S = (110)\) is extracted. It means that the error position is the 6-th from left.

2.2 Hamming coding+1 data hiding (H1DH)

For H1DH, \(k\) secret bits \(\delta\) embed to \(y_b = (y_1, \ldots, y_n)\) of \(n (= 2^k − 1)\) pixels using \(COV(1, n, k)\) by flipping at most one bit. An extra 1-bit embeds at \(2^k\) by using OPAP. That is, first \(k\) was embedded by using syndrome \(\hat{y} = y_b \oplus e(\delta - Hy_b)\). If \(\hat{y} \cdot H^T = \delta\), there is no needed any action. For embedding \(\delta\), one change is needed with probability \((2^k - 1)/2^k\). To embed \(k + 1\) bits, they append one pixel like \(y = (y_1, \ldots, y_n, y_{n+1})\). By \(y_i \pm 1\), its \(b(y_i)\)
becomes the same binary $b(y_i) \oplus 1$, while $2b(y_i)$ can either be 0 or 1 after $b(y_i) \oplus 1$, where $b(\cdot) = (y_i \mod 2)$ and $2b(\cdot) = ([y_i/2] \mod 2)$ are LSB and 2LSB, respectively.

$$\begin{align*}
(\delta_1, \ldots, \delta_k)^T &= H \cdot y_b \\
\delta_{k+1} &= (2b(y_1) + \cdots + 2b(y_{2^k-1}) + b(y_{2^k})) \mod 2
\end{align*}$$

There are four embedding optimal rules for $H1DH$.

- **O1**: When Eqs. (3) and (4) are satisfied, no action.
- **O2**: When only Eq. (3) is fulfilled, process $y_{2^k} \pm 1$.
- **O3**: When only Eq. (4) is satisfied, process $y_i \pm 1$.
- **O4**: When both of two are not satisfied, process $y_{2^k} \pm 1$.

### 2.3 Hamming+$k$ scheme with $m$ overlapped pixels ($H_kmDH$)

The $H_kmDH$ [19] is the extended method of $H_kDH$ [18]. Suppose that one block is composed of $(2^k-1-m-2)$ pixels $(y_1, \ldots, y_{2^k-1-m-2})$ with $m$ overlapped pixels.

![Fig. 1. Block diagram of $H_kmDH$ with $k=3$ and $m=3$.](image)

$H_kmDH$ uses $COV(1, 2^k-1, k)$ to embed $k$ secret bits $(\delta_1, \ldots, \delta_k)$ into the first $(2^k-1)$ pixels, and then embed the other $k$ secret bits into the last $(2^k - 1)$ pixels. The first $(2^k - 1)$ pixels are $(y_1, \ldots, y_{2^k-1})$, and the other pixels are $(y_{2^k-m}, \ldots, y_{2^k-m-2})$. As shown in Fig. 1, consider the example $k=3$ (i.e., using $(7, 4)$ Hamming code) and $m=3$. There are total 11 $(=2^3-3-2)$ pixels with 3 overlapped pixels. To embed secret bits, we apply Eqs. (5) and (6) to each pixels $(2^k-1)$.

$$\begin{align*}
p &= ((b(y_1), \ldots, b(y_{2^k-m-1}), b(y_{2^k-m})\oplus 2b(y_{2^k-m}), \ldots, b(y_{2^k-1})\oplus 2b(y_{2^k-1}))) \\
\quad (\delta_1, \ldots, \delta_k)^T &= H \cdot p^T
\end{align*}$$

$$\begin{align*}
q &= ((2b(y_{2^k-1})\oplus 2b(y_{2^k-1}), (y_{2^k}) \ldots, b(y_{2^k+1-m-2}))) \\
\quad (\delta_{k+1}, \ldots, \delta_{2k})^T &= H \cdot q^T
\end{align*}$$

We embed $k$ secret bits into the first seven bits $p$ for Eq. (5), and then embed the other $k$ secret bits into the other seven bits $q$ of Eq. (6). Except the $m$ overlapped pixels, other pixels adopt the LSBs. The overlapped pixels in Eq. (5) (respectively, Eq. (6)), are the XOR-ed
results of 1st LSB and 2nd LSB, and 2nd LSB, respectively.

The embedding algorithm is as follows.

| The embedding algorithm \((H_k_mDH), m=3\) |
|------------------------------------------|
| **Input:** Original image \(O_I\)  |
| **Output:** Stego image \(S_I\)  |
| (Step 1) Read one block from the cover object, and generate a codeword, \(p\) by Eq.(5). |
| (Step 2) \(S_1 = H \cdot p^T\) and \(S'_1 = S_1 \oplus \delta^{j+2}, j = 1\). |
| (Step 3) If \((S'_1 \leq (2^k - m), \pm(x_{S'_1})\), else if \((S'_1 \leq (2^k - 1), when (2b)b(x_{S'_1}) = ([00] \ or \ [10])); y_{S'_1} + 1\) and when \((b(y_{S'_1}) = ([01] \ or \ [11])); y_{S'_1} - 1\). |
| (Step 4) Generate another codeword, \(q\) by Eq.(6) |
| (Step 5) \(S_2 = H \cdot q^T\) and \(S'_2 = S_2 \oplus \delta^{j+2}\). |
| (Step 6) If \((S'_1=S_2), \pm(\pm_{y_{2k+1-m-2}}), goto (Step 5)\) |
| else if \((S'_1 \neq S_2), \{\) |
| If \(S'_2 \geq 2^k, \pm(x_{S'_2})\), else if \((S'_2 \geq 2^k - m \ \& \ S'_2 \leq 2^k + 1), when (2b)b(x_{S'_2}) = ([00] \ or \ [10])); x_{S'_2} - 1\) and when \((b(y_{S'_2}) = ([01] \ or \ [11])); x_{S'_2} + 1\). |
| \} |
| (Step 7) Go to Step 1 until not end of block. |

3. The Proposed Schemes

3.1 Design concept

All Hamming code based DHs are shown in Fig. 2. The ERs for HDH, H1DH, \(H_k\_mDH\), and \(H_k\_mDH\) are \(k/n, (k+1)/n, 2k/n, 2k/(2n-m)\), respectively (see Figs. 2(a)~(d)). As shown in Fig. 2(e), the proposed SH\(_k\_mDH\) is to sequentially use \(H_k\_mDH\) to embed secret bits. The ER of the proposed SH\(_k\_mDH\) is derived as Eq. (7).

\[
ER = \frac{\alpha k}{\alpha(n-m)+m} = \frac{k}{(n-m)+m/\alpha} = \frac{k}{n-m} \quad (\text{for } \alpha \rightarrow \infty)
\]  

(7)

The value of \(k/(n-m)\) is larger than \(k/(2n-m)\) of \(H_k\_mDH\). For instance, when \(\alpha = 3, n = 7, \text{ and } m = 3\), \(ER = 9/15 = 0.75\). In aspect of PSNR, the SH\(_k\_mDH\) is the same to \(H_k\_mDH\). We can reduce the distortion by using OPAP and LSB.

In our proposed block overlapping approach, an image is divided by blocks of sized \((1 \times n)\) pixels. At this time, every block has \(2m\) overlapped pixels except for the first block including \(m\) overlapped pixels. The main difference between the SH\(_k\_mDH\) and the \(H_k\_mDH\) is that we sequentially embed secret bits with \((n-2m) \geq 0\) (see Fig. 2(e)).
3.2 Embedding Procedure

For brief explanation, we here assume that $k = 3$ (i.e., by COV (1, 7, 3) covering function) and $m = 3$ for the proposed SH$k_m$DH, which can embed 3 secret bits into 7 overlapped cover pixels.

The detailed procedure of the embedding is as follow.

Fig. 2. Diagrammatical representation of Hamming code based DH:
(a) HDH (b) H1DH (c) H$k$DH (d) H$k_m$DH (e) the SH$k_m$DH.
The embedding algorithm (SH$_k$-mDH)

**Input:** Original image OI

**Output:** Stego image SI

1. **Read one block, overlapped ($2^k - 1$) pixels ($x_i^{2^k-1}$), from the cover object and generate a codeword,** $p \in F_2^n$, by Eq.(8).

2. **Calculate the syndrome** $S = H \cdot p^T$ and $S' = S \oplus (d^j_{j+2})$.

3. **if $S \neq 0$**
   - if ($S \leq 4$) figure out Eq. (9),
   - else if ($S \leq 7$) figure out Eq. (10).

4. **if not end of the block,** $i = i + 4$, $j = j + 3$ and go to Step 1.

$$p = (b(x_1) \oplus b\left(\left\lfloor \frac{x_1}{2} \right\rfloor\right), \ldots, b(x_4) \oplus b\left(\left\lfloor \frac{x_4}{2} \right\rfloor\right), b\left(\left\lfloor \frac{x_5}{2} \right\rfloor\right), \ldots, b\left(\left\lfloor \frac{x_7}{2} \right\rfloor\right)) \quad (8)$$

$$x_{i'}' = \begin{cases} x_{i'}' + 1 & \text{if } (2b(x_{i'}') = [00] \text{ or [10])} \\ x_{i'}' - 1 & \text{if } (2b(x_{i'}') = [01] \text{ or [11])} \end{cases} \quad (9)$$

$$x_{i'}' = \begin{cases} x_{i'}' + 1 & \text{if } (2b(x_{i'}') = [01] \text{ or [11])} \\ x_{i'}' - 1 & \text{if } (2b(x_{i'}') = [00] \text{ or [10])} \end{cases} \quad (10)$$

**Fig. 3.** Embedding example using SH$_k$-mDH.
Example 1: We demonstrate the embedding procedure using given \( x \) and \( \delta \) as shown in Fig. 3. For embedding, first, \( x = (162, 162, 162, 162, 162, 164, 161) \) and \( \delta = (010) \) is given. The codeword \( p \) is obtained from a vector \( x \) using Eq. (8) and then figure out syndrome \( S = H p = (001) \) using the codeword \( p \). First 4 bits of the codeword is obtained by using 2LSB \( \oplus \) LSB. The other 3-bits are extracted from 2LSB of the \( x \) directly. \( S(= S \oplus \delta) \) means the codeword \( p \) has new error \( e = (0010000) \). If we flip the third bit, then syndrome of the codeword \( p' \) is equal to \( S = (011) \). However, it is not acceptable to modify a bit in the codeword \( p \) for embedding secret bits directly, because the codeword \( p \) cannot store any bits. Therefore, we must modify a pixel in the vector \( x \) directly according to the rule in Eqs. (9) and (10). In this case, since \( S' \) is 3, it is belonging to the rule of Eq. (9). Thus, because \( b2|b(x_{S'}) \) is equal to “[10],” figure out \( x_{S'} = x_{S'} + 1 \). That is, \( x_3 \) become 163.

For second, \( x = (162, 164, 161, 158, 156, 156, 157) \), including three overlapped pixels \((162, 164, 161)\) of former pixel, and secret bits \( \delta = (001) \) are given. First, we generate the codeword \( p \) from the vector \( x \) using Eq. (8) and calculate both \( S \) and \( S' \) in the same way in the first case. Since \( S' \) is less than or equal to 7, it is belong to the rule of Eq. (10). That is, \( x_e = 155 \). For third and fourth cases, we can get the same way like Fig. 3 (a) and (b).

3.3 Theoretical Estimation of PSNR

The AMSE of SH\(_k\)\_mDH is estimated as follows. Consider a codeword \( C_i \) of length \( n (=2^k-1) \), where we could embed \( k \) secret bits. Because the proposed SH\(_k\)\_mDH embeds every \( k \) secret bits with \( m \) overlapped pixels recursively. Therefore, embedding secret bits in the previous codeword \( C_{i-1} \) and the next codeword \( C_{i+1} \) may modify the pixels in the current codeword \( C_i \).

Originally, when the modification places in overlapped area of are the same, we should modify other two modification places to avoid the collision. There are \( j \) modifications, \( 0 \leq j \leq 3 \), in \( C_i \).

Case (1) \( j=0 \):

\[
AMSE_0 = \frac{(n-m+1)(\frac{1}{n+1} \times \frac{n-m+1}{n+1}) \times 0^2}{n}.
\]

Case (2) \( j=1 \):

\[
AMSE_1 = \frac{(n-m+1)(\frac{1}{n+1} \times \frac{n}{n+1} \times \frac{n-m+1}{n+1}) \times 1^2}{n} + \frac{(m \times \frac{1}{n+1} \times \frac{n-m+1}{n+1}) \times 1^2}{n}.
\]

Case (3) \( j=2 \):

\[
AMSE_2 = \frac{(n-m+1)(\frac{1}{n+1} \times \frac{m}{n+1} \times \frac{n-m+1}{n+1}) \times 1^2}{n}.
\]
Finally, the AMSE of $SH_k_{m DH}$ is derived as follows. For simplicity, we may neglect the cases that the collision in overlapped area to get a simple form of AMSE as $\frac{1}{2k} + \frac{m}{2^{k-1}}$ (see Eq.(11)).

$$AMSE_2 = \left( \frac{m}{n+1} \times \frac{1}{n+1} \times \frac{m}{n+1} \right) \times (I^2 + I^2) / n + \left( \frac{n-m+1}{n+1} \times \frac{1}{n+1} \times \frac{m}{n+1} \right) \times (I^2 + I^2) / n.$$  

Case (4) $j=3$:

$$AMSE_3 = \left( \frac{m}{n+1} \times \frac{n}{n+1} \times \frac{m}{n+1} \right) \times (I^2 + I^2) / n + 2 \times \left( \frac{m}{n+1} \times \frac{1}{n+1} \times \frac{m}{n+1} \right) \times I^2 / n.$$  

Finally, the AMSE of $SH_k_{m DH}$ is derived as follows. For simplicity, we may neglect the cases that the collision in overlapped area to get a simple form of AMSE as $\frac{1}{2k} + \frac{m}{2^{k-1}}$ (see Eq.(11)).

4. Experiment and Comparison

4.1 Experimental Results

To demonstrate the performance of the proposed method, we compare it with HDH, H1DH, Kim & Yang’s DH, and Hk mDH by using PSNR and EC. In this experiment, we use 512×512 original grayscale images as cover images [19]. The PSNR is an expression for the ratio between the maximum possible value (energy) of a signal and the energy of distorting noise that affects the quality of its representation. The visual quality of digital images is very subjective because of personal bias. This is why we should establish quantitative/empirical measures to compare the effects of image quality. Using the same set of tests images, different DH algorithms can be compared fairly to identify which algorithm has better results. If a stego image can closely resemble the original, then it is a good algorithm. In this aspect, the PSNR may help us to do this.

The value of MSE used to measure average the squared intensity differences between a distorted image and reference image. If an image has a small MSE value, it has a relatively good visual quality. The MSE between two image $x$ and $y$ is

$$MSE(x, y) = \frac{1}{N} \sum_{i=1}^{N} (x_i - y_i)^2 .$$  

The value $e_i = x_i - y_i$ denotes the difference between the original and distorted signal. The MSE is a function for PSNR measure where $L$ is the dynamic range of allowable pixel intensities. For example, for an 8-bit per pixel image, $L = 2^8 - 1 = 255$.  

$$\frac{1}{L} \sum_{i=1}^{L} (x_i - y_i)^2 .$$
PSNR = 10\log_{10} \frac{L^2}{MSE} \hspace{1cm} (13)

Meanwhile, the EC adopts to compare the performance of the DH through measuring embedding capacity of the cover image. Obviously, most DHs endeavor to increase EC without degrading the quality of stego image. Eq. (14) implies that EC is the ratio of the number of message bits (||δ||)to the total number of pixels.

\rho = \frac{||δ||}{N \times N} \hspace{1cm} (14)

**Table 1** shows the comparison among HDH, H1DH, Kim and Yang’s DH, H_k mDH, and the proposed SH_k mDH which are HC based DHs. Each ER of those is 0.43, 0.5, 0.54, 0.54, and 0.75, respectively. The ER of SH_k mDH is higher than those of other DHs. Moreover, the SH_k mDH is also better than that of Kim and Yang’s DH in the aspect of ER and PSNR. The PSNR of H_k mDH is slightly higher than that of SH_k mDH, but ER of the SH_k mDH is 0.75 and so it is higher than 0.54 of H_k mDH.

**Table 1. Performance comparison of all Hamming-like DHs**

| Images  | HDH   | H1DH  | Kim and Yang’s DH | H_k mDH | SH_k mDH  |
|---------|-------|-------|-------------------|---------|-----------|
|         | PSNR  | \rho  | PSNR  | \rho  | PSNR  | \rho  | PSNR  | \rho  | PSNR  | \rho  | PSNR  | \rho  |
| Baboon  | 0.43  | 58.5731 | 0.5  | 53.96 | 0.54  | 56.0045 | 0.54  | 54.2749 | 0.75  |
| Barbara | 0.43  | 58.5851 | 0.5  | 53.93 | 0.54  | 56.0121 | 0.54  | 54.2874 | 0.75  |
| Boats   | 0.43  | 58.6066 | 0.5  | 53.94 | 0.54  | 55.9930 | 0.54  | 54.2823 | 0.75  |
| Goldhill| 0.43  | 58.5937 | 0.5  | 53.95 | 0.54  | 56.0005 | 0.54  | 54.2974 | 0.75  |
| Airplane| 0.43  | 58.5755 | 0.5  | 53.95 | 0.54  | 55.9985 | 0.54  | 54.2943 | 0.75  |
| Lena    | 0.43  | 58.5772 | 0.5  | 53.95 | 0.54  | 56.0164 | 0.54  | 54.3018 | 0.75  |
| Peppers | 0.43  | 58.5616 | 0.5  | 53.95 | 0.54  | 56.0063 | 0.54  | 54.2832 | 0.75  |
| Tiffany | 0.43  | 58.6239 | 0.5  | 53.96 | 0.54  | 56.0104 | 0.54  | 54.2949 | 0.75  |
| Zelda   | 0.43  | 58.5806 | 0.5  | 53.99 | 0.54  | 56.0112 | 0.54  | 54.2902 | 0.75  |
| Average | 0.43  | 58.5863 | 0.5  | 53.95 | 0.54  | 56.0058 | 0.54  | 54.2896 | 0.75  |

**Table 2. Performance of SH_k mDH for 1 ≤ m ≤ 3**

| Images  | m = 1 | m = 2 | m = 3 |
|---------|-------|-------|-------|
|         | PSNR  | \rho  | PSNR  | \rho  | PSNR  | \rho  |
| Baboon  | 56.3327 | 0.5  | 55.4245 | 0.6  | 54.2749 | 0.75  |
| Barbara | 56.3379 | 0.5  | 55.4020 | 0.6  | 54.2874 | 0.75  |
| Boats   | 56.3391 | 0.5  | 55.4134 | 0.6  | 54.2823 | 0.75  |
| Goldhill| 56.3486 | 0.5  | 55.4075 | 0.6  | 54.2974 | 0.75  |
| Airplane| 56.3399 | 0.5  | 55.3981 | 0.6  | 54.2943 | 0.75  |
| Lena    | 56.3422 | 0.5  | 55.3820 | 0.6  | 54.3018 | 0.75  |
| Peppers | 56.3403 | 0.5  | 55.4063 | 0.6  | 54.2832 | 0.75  |
| Tiffany | 56.3469 | 0.5  | 55.4177 | 0.6  | 54.2949 | 0.75  |
| Zelda   | 56.3305 | 0.5  | 55.3951 | 0.6  | 54.2902 | 0.75  |
| Average | 56.3397 | 0.5  | 55.4051 | 0.6  | 54.2896 | 0.75  |
H1DH is superior in quality compared to other DHs, but the proposed SH$k_m$DH is the highest among all DHs in the respect of PSNR. Especially, SH$k_m$DH and H$k_m$DH will be able to determine ER using the parameter $m$. In fact, there is a trade-off between PSNR and ER. We may adjust the value of $m$ according to our application.

Table 2 shows the performance of SH$k_m$DH based on variable $m$. For $1 \leq m \leq 3$, ERs are 0.5, 0.6, and 0.75, respectively. That is to say, ERs increase as increasing the parameter $m$. Meanwhile, PSNRs decrease. The strength of our SH$k_m$DH is that users may embed a large number of bits by increasing the value of $m$ with retaining the high PSNR.

Table 3. Performance comparison of proposed scheme and previous schemes (when ER = 0.34)

| Images     | HDH ($n=7, k=3$) | H1DH ($n=7, k=3$) | Kim and Yang’s DH ($n=7, k=3$) | H$k_m$DH ($n=7, k=3$) | SH$k_m$DH ($n=15, k=4$) |
|------------|------------------|-------------------|-------------------------------|------------------------|------------------------|
|            | PSNR             | PSNR              | PSNR                          | PSNR                   | PSNR                   |
| Baboon     | 58.1147          | 60.2717           | 57.8956                       | 57.9809                | 58.9346                |
| Barbara    | 58.1150          | 60.2090           | 57.8783                       | 57.9931                | 58.9332                |
| Boats      | 58.1231          | 60.2636           | 57.8873                       | 57.9963                | 58.9249                |
| Goldhill   | 58.1094          | 60.2299           | 57.8964                       | 58.0172                | 58.9320                |
| Airplane   | 58.1130          | 60.2636           | 57.8823                       | 58.0117                | 58.9382                |
| Lena       | 58.1241          | 60.2412           | 57.8967                       | 58.0127                | 58.9290                |
| Peppers    | 58.1295          | 60.2439           | 57.9062                       | 58.0261                | 58.9316                |
| Tiffany    | 58.1819          | 60.2579           | 57.9517                       | 58.0209                | 58.9558                |
| Zelda      | 58.1186          | 60.2642           | 57.8952                       | 58.0008                | 58.9263                |
| Average    | 58.1254          | 60.2494           | 57.8988                       | 58.0066                | 58.9339                |

Fig. 4. Comparison of Goldhill images generated from various schemes: (a) original (b) HDH (58.12 dB), (c) H1DH (60.24 dB), (d) Kim and Yang’s DH (57.89 dB), (e) H$k_m$DH (58.01 dB), and (f) SH$k_m$DH (58.92 dB).
Table 3 shows the comparison of PSNR between proposed SH\(_k_m\)DH and previous schemes such as HDH, H1DH, Kim & Yang’s DH, and H\(_k_m\)DH when ER is 0.34. As shown in Table 3, even if the same amount of data are hidden in given images, the PSNR of the stego images appears different results depending on the performance of the algorithm. In this experiment, SH\(_k_m\)DH (in case of COV (15, 4)) had higher PSNRs than HDH, Kim and Yang’s DH, and H\(_k_m\)DH, while PSNR of H1DH is 0.13 dB higher than that of SH\(_k_m\)DH. In fact, SH\(_k_m\)DH maximize ER while maintaining PSNR.

When ER = 0.34, Fig. 4 represents the comparison among original images and stego images such as (a) original (b) HDH, (c) H1DH, (d) Kim and Yang’s DH, (e) H\(_k_m\)DH, and (f) SH\(_k_m\)DH. The PSNR of H1DH is the highest. The PSNR of SH\(_k_m\)DH is higher 0.92 dB than that of H\(_k_m\)DH.

As shown in Fig. 5, experiments were conducted to verify the performance of the proposed method for 1\(\leq m\leq 3\) and the ER was (0.1…0.75). The grayscale Lena image is used for the experiment. For the same ER, the PSNR of \(m=1\) is higher than those of \(m=2\) and 3.

Entropy [21] is a statistical measure of randomness that can be used to characterize the texture of the input image. The normalized histogram is an estimate of the underlying probability of pixel intensities, \(i.e., N/h(i)\), where \(h(i)\) denotes the histogram entry of intensity value in an image and \(N\) is the total number of pixel. The entropy of an image is computed as:

\[
E = \sum_i h(i) \log \frac{N}{h(i)}
\]  

(15)

When the relative entropy \(E\) lying between two probability distribution functions is zero, the system is perfectly secure. In Table 4, the entropy of original image (Entropy O), stego image (Entropy S), and the difference between Entropy O and Entropy S are enumerated. It appears that when the number of bits in the secret message increases, the relative entropy in stego
image also increases. Because the differential entropy is approximated to zero, it means the proposed scheme endures an attack of steganalysis tools.

Table 4. Comparison of relative entropy between original images and stego images.

| Image     | ER | Entropy O | Entropy S | Difference |
|-----------|----|-----------|-----------|------------|
| Lena      | 0.3| 7.3938    | 7.3937    | 0.0001     |
| Baboon    | 0.3| 7.3579    | 7.3579    | 0.0000     |
| Boats     | 0.3| 7.1237    | 7.1256    | 0.0018     |
| Goldhill  | 0.3| 7.4777    | 7.4794    | 0.0016     |
| Tiffany   | 0.3| 6.6055    | 6.6064    | 0.0008     |

RS steganalysis [22] was developed with the intent to detect embedded secret messages using LSB replacement. For RS steganalysis, the discrimination function (DF) \( f \) was used (Eq.(16)), which is to capture the smoothness or “regularity” of the group of pixels.

\[
f(x_1, x_2, \ldots, x_n) = \sum_{i=1}^{n-1} |x_{i+1} - x_i|
\]  

(16)

A pixel groups can be classified as three types: R, S, and U.

\begin{itemize}
  \item Regular groups: \( G \in R \iff f(F(G)) > f(G) \).
  \item Singular groups: \( G \in S \iff f(F(G)) < f(G) \).
  \item Regular groups: \( G \in U \iff f(F(G)) = f(G) \).
\end{itemize}

where \( F(G) = (F_M(x_1), F_M(x_2), \ldots, F_M(x_n)) \). Thus, total number of R groups will be increase than the total number of S groups. When the parameters satisfy Eq.(17), it indicates that there is no hidden data in the respective image. When an image has hidden data, \( R_M \) and \( S_M \) increases, whereas \( R_M \) and \( S_M \) decrease and detected by RS steganalysis.

\[
R_M \approx R_{-M} \text{ and } S_M \approx S_{-M}
\]

(17)

According to increase ER, the difference between \( R_M \) and \( S_M \) is to zero. After flipping the LSB of about 43% of pixels, it become \( R_M \neq S_M \) when \( M = [1 \ 1 \ 0 \ 0 \ 1] \). As shown in Fig. 6 (a), it is a few to find the \( R_M \approx R_{-M} \) and \( S_M \approx S_{-M} \). Meanwhile, in Fig. 6 (b),(c), and (d), we can see \( R_M \approx R_{-M} \) and \( S_M \approx S_{-M} \) until about 100%. Therefore, HDH, \( Hk_{MDH} \), and \( SHk_{mDH} \) have safety zones avoiding steganalysis detection.
5. Conclusion

$H_{k \cdot m}DH$ is better than previous HC based DHs in the aspect of ER. In this paper, the reason why we proposed $SH_{k \cdot m}DH$ was that $H_{k \cdot m}DH$ is not applied sequentially. As a result, could the sequential approach increases the ER to 0.75 when $m = 3$. In $SH_{k \cdot m}DH$, a collision may occur when the same positions in two overlapping blocks need to be corrected at the same time. In this paper, we solve the collision problem using OPAP and LSB. In addition, we demonstrated the superiority of $SH_{k \cdot m}DH$ via experiments and proofs.

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