An energy efficient optimized cluster establishment methodology for sensor nodes in WSN

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ABSTRACT

The compatibility of WSN is with various applications such as; healthcare and environmental monitoring. Whereas nodes present in that network have limited ‘battery-life’ that cause difficulty to replace and recharge those batteries after deployment. Energy efficiency is a major problem in the present situation. In present, many algorithms based on energy efficiency have been introduced to improvise the conservation of energy in WSN. The LEACH algorithm improvises the network lifetime in comparison to direct transmission and multi-hop, but it has several limitations. The selection of CHs can be randomly done that doesn’t confirm the optimal solution, proper distribution and it lacks during complete network management. The centralized EE optimized cluster establishment approach (OCEA) for sensor nodes is proposed to decrease the average energy dissipation and provide significant improvement. The proposed EE WSN model with the sensor nodes is examined under a real-time scenario and it is compared with state-of-art techniques where it balances the energy consumption of the network and decreasing the cluster head number.

Keywords: Base station, Cluster head, Energy efficient, LEACH, EE, Low energy adaptive clustering hierarchy, Wireless sensor network

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1. INTRODUCTION

WSN is a real-time environment and it usually contains several types of user devices. The main aim of these mechanisms is to sense the data from the deployed environment and transfer towards BS [1]. WSN discover broad utilization in many fields like military applications, environmental monitoring, target tracking and structural health monitoring. WSN is composed of spatially dispersed sensors that monitor the physical parameters and environmental [2]. It interacts and gathers the data to the sink node for the analysis of the end-user. Transmitting data packets from nodes to the base station can be accomplished by various routing techniques. Because of limited nodes of power supplies, these types of components must be taken the benefits as much as they can [3]. The clustering approach comes out to be a very effective and most promising method to improve the EE [4]. Past few years, many algorithms like EE based clustering routing have been introduced to improvise the conservation of energy in WSN [3-5].

Recently, many types of research have come up with various ideas in order to minimize the energy and improvise the network longevity for proper resource utilization. In this process, the routing algorithm plays a vital role. Clustering constructs hierarchy clusters or sets of the sensing nodes, which gathers and transmits data to its CH. The CH groups the data and transmits to BS/sink that acts as middleware among
network and end-user. LEACH and clustering algorithm are the classical protocol that considers the energy for hierarchical data routing [6]. The network can be grouped onto the clusters and SN (sensor node) which sends their data to corresponding CH. The protocol randomly chooses CHs in a stochastic manner for each round. CH interacts with every single cluster node which is known as the SN to sense and collect the data.

CH allocates schedules of time division multiple access (TDMA) to its corresponding cluster member (CM). The member node can transfer the data during the allocated time-slot. The data can be checked for the redundancy and compacted before interacting with the sink-node. The CHS directly interacts with BS in the protocol of LEACH; hence the consumption of power is transmitting the data from CH to BS that will be compared to communication among CHs. As an outcome, CHs exhausted its energy within a short time. On the other side, in order to overcome this issue, multi-hop communication can be very helpful, but still, it is not effective in small networks.

Selecting CH is the sophisticated job as different factors that have to be assumed for the best selection node in a cluster [7]. These different factors consist of distance among mobility, throughput, residual energy, and nodes of every single node. The algorithm of LEACH improvises the network lifetime in comparison to multi-hop and direct transmission but still, it has several limitations. The selection of CHs can be randomly done that doesn’t confirm the optimal solution and proper distribution. The nodes with less energy have the same priority as with high energy levels to be selected as the CH. While, lower residual energy node achieves chosen to serve as the CH, and quickly expires the resulting in a short network span [8].

LEACH-ME [9] chooses the SN with a lower mobility factor to take benefits on the role of CH. The CBR (cluster-based routing) protocol for MNs (mobile nodes) [10] uses the design of a cross-layer to manage mobility. The MBC (mobility based clustering) [11] protocol selects SN to CH on the basis of its remaining mobility and energy. The main drawback of this protocol is the threshold function that choosing the SN to be CH that may be very smaller than the LEACH, because it is multiplied simultaneously by the mobility factor and energy factor that leads to unbalance CHs number, thereby maximizing the dissipation of energy. Moreover, LEACH-ME [12] utilizes the fuzzy logic to combine more than one factor for the selection of CH. In this paper, the centralized EE clustering routing protocol for sensor nodes is introduced to reduce the average dissipation of energy and provide promising improvement. The main contribution of this work: the optimized Clustering algorithm is improved that periodically chooses the optimal set of the CHs according to an average node of hybrid energy and average speed of node. This is crucially eminent that it utilizes the distributed clustering protocol in order to choose the CHs on the basis of current node speed and residual node energy.

2. LITERATURE SURVEY

In the past few decades, several studies have been devoted to optimizing the cluster size in the WSNs. The main idea of dynamic clustering is introduced HEED [13], LEACH [14] aims to improvise EE. Anyways, performing the dynamic clustering executes more overhead on the network. This problem has been not solved in the model of unequal-clustering-size (UCS) [15] as the first scheme of unequal clustering. The main target of schemes [16, 17] is to address the problem oh hot-spot that could happen in the clusters around BS since these types of clusters need extra power to relay data-packets from the outer clusters. In UCS, the CHs consumption of balanced energy can be improved by adjusting the number of nodes in various clusters based on traffic load.

Thus, UCS improvise EE that outcomes in the lifetime of a longer network. Additionally, utilizing a 2-hop of inter-cluster communication that leads to reduce the transmission distance in comparison to the LEAH. Anyways, the network is parted into 2 layers whereas UCS is inapplicable for the larger size of networks. Furthermore, the energy consumption of inter-cluster is not assumed in the UCS. The randomization technique can be distributed energy between SNs in the network. LEACH [18, 19] considers each node that has the radio powerful to reach directly to the sink node. But operating radio in its maximum power that depletes the battery rapidly.

Additionally, it can be observed that the consumption of balanced energy is accomplished by assigning the larger size of the cluster to CHs that belonging to the outer layers. It has very little communication [20], whereas this problem is not considered in the energy-balanced-data-gathering (EBDG) [21]. To improve the lifetime of the network, the coverage time optimization-based method is represented [22] that intends to compute the optimal size of the cluster in various layers. Anyways, one of the main desires of such a method is several clusters; however, it is not feasible [23]. Arranging the ACT (Transmission ranges) and cluster sizes methods have been represented by [23] that intends to overcome the problem of a hot spot and improve the network lifetime. In the ACT, the network topology contains many layers and cluster size that belonging to every layer is defined according to distance from the BS. Therefore,
this property can resolve the problem of an energy hole. Anyways, clusters size fitting is very larger to the outermost clusters that lead to improvising the energy consumption of intra-cluster as outer layers. Additionally, these new locations are chosen as the CH deviates from ideal ones and this property creates very energy distribution.

In [24] park method defines CH selection based on the K-Means algorithm. In the network, the SNs are gathered into the clusters. Before cluster formation, each CH is chosen randomly based on the minimum distance of Euclidean among MNs and centroid points. The main problem with the park method is the time that takes for the formation of a cluster. The EBCAG based scheme (energy-balancing-unequal-clustering-approach for gradient) is represented by [25]. The EBCAG intends to improvise the size of the cluster in various network layers to avoid unbalanced nodes of energy consumption that belonging to the various layers. The cluster size is computed by the help of gradient values of CHs that are calculated based on received data from given nodes with the gradient values. This method balances the CHs power consumption as well as it disables the disadvantages of optimization-based cluster size methods.

Anyways, EBCAG creates a smaller size of the cluster and it is the opposite of being the network scalability. In DBS [26], the weal point has been removed by giving the mathematical framework to improvise the size of the cluster in various network layers. The DBS distributes network into the coronas. After that, various clustering methods are implemented in every single corona that outcomes in improved energy conservation and LB (load balancing). The algorithm of sub-clustering named OCCN (optimal-clustering in circular-networks) is introduced by [27] to improve the lifetime of network and minimize dissipated energy. In the end, consumption of energy can be improved by dving the network into the same size of clusters in a distributed manner. Anyways, the same size of cluster leads to a difference in the consumption of energy during inter-cluster communication that outcomes at initial network partitioning.

3. RESEARCH METHOD

A WSN model with the mobile nodes (MNs) is examined in this paper. It is assumed that this model of the network consists MNs that uniformly scattered in the area (z × z meters) and WSN clustered diagram is shown in Figure 1. The assumptions of the network model are given below:

a) The BS is fixed and located in a sensing field sensor.
b) The FNs and BS are nor doing the energy constrain, while the MNs are the energy-constrained.
c) Each node knows its velocity, current position, and energy level.
d) Each node can use the power setting to change the number of transmitted power via RSSI (Received-Signal-Strength-Indication).

Figure 1. WSN Clustered Diagram

These above-given assumptions as reasonable for several real networks. Such as FNs, can be powered by the help of energy harvesting and have a continuous supply of power [28]. In this case, each
node knows its current speed and position that can be obtained by location and GPS algorithms [29, 30]. Similarly, each node can evaluate its level; of energy according to specific hardware of the energy model.

The MNs energy consumption is measured by similar model as in [31], where the distance between the receiver and transmitter is considered. If a distance is lesser than a threshold then the open space channel-model is considered with energy loss (θ) or else the model of multipath channel (μ) is considered. So in order to send and receive b number of bit message signal at distance, the cost of transmitting and receiving are given as:

\[ D_T(b,a) = b D_{ed} + bc_{sc} \theta \] (1)

Where a distance is lesser then \( a_0 \)

\[ D_T(b,a) = b D_{ed} + bc_{mp} \mu \] (2)

Where a distance is greater or equal to \( a_0 \)

\[ D_R(l) = b D_{ed} \] (3)

Here, \( D_{ed} \) denotes for energy dissipation/bit at process of transmitting and receiving, \( c_{sc} \) denotes for parameter of amplify energy and corresponds to open space channel-model. Moreover, \( c_{mp} \) is a parameter of amplify energy and corresponds to multipath channel-model.

All MNs from a network is given by \( E = \{ e_1, e_2, e_3, \ldots, e_n \} \). Starting phase, a round \( g \) for each \( e_i \) transmit message towards the base station, where \( i \) values ranges from 1 to \( n \). However, the transmitted message contains a node location, \( H(e_i) \) present node speed and \( D(e_i) \) energy level of node.

In accordance to this transmitted message, initially base station computes the average of node energy follow as:

\[ \bar{D}_n = \frac{1}{n} \sum_{i=1}^{n} D(e_i) \] (4)

Also compute the average of node speed follow as;

\[ \bar{H}_n = \frac{1}{n} \sum_{i=1}^{n} H(e_i) \] (5)

Where, \( e_i \in E \). The base station chooses a number of nodes \( K \) from \( E \) set (\( K \subseteq E \)) and given as;

\[ K = \begin{cases} \{ K_1 \cap K_2 \}, & \text{if } |K_1 \cap K_2| \geq L_b \\ \{ K_1 \cup K_2 \}, & \text{if } |K_1 \cap K_2| < L_b \end{cases} \] (6)

The symbol of \( L_b \) denotes for optimal channel heads (CHs) number [31], the \( K_1 \) and \( K_2 \) shows set of nodes, and \( |K_1 \cap K_2| \) shows the cardinality of \( K_1 \cap K_2 \). Moreover, energy level of nodes is above then the threshold/average energy value, and the computed current speeds below then the value of average speed is:

\[ K_1 = e_j | D(e_j) \geq \bar{D}_n, \forall e_j \in E \] (7)

\[ K_2 = e_j | H(e_j) \geq \bar{H}_n, \forall e_j \in E \] (8)

While considering a condition where \( |K_1 \cap K_2| \geq L_b \), this implies that set of nodes in \( K \) are having low speeds and high energy levels are qualified to become CHs during that particular round. Whereas considering a condition \( |K_1 \cap K_2| < L_b \), this implies that set of nodes in \( K \) are having low speeds or high energy levels are probable to become CHs during that particular round.

The set of nodes from \( K \) is considered to perform an optimization process to get optimal CHs that becomes CHs at a particular round. Here several processing levels is used to get optimal CHs such as; initialization level, neighborhood finding level and final/completion level.

At initialization level, initially stochastically set of CHs is generated and denoted by \( L = \{ l_1, l_2, l_3, \ldots \} \) from \( K \) and \( m = 0 \) in order to ensure \( |L| = L_b \). At neighborhood finding level, \( m = m + 1 \) is considered to find out a novel set of nodes \( N \), which shows random perturbation of \( L \) nodes. Therefore, the new \( \bar{a} \) and \( \bar{b} \) coordinates of nodes in \( N \) are computed through \( \alpha \) and \( \beta \) coordinates of nodes in \( L \).
\[ \tilde{a} = \alpha + \text{rand}(-c_{\text{max}}, c_{\text{max}}) \]  
\[ \tilde{\beta} = \beta + \text{rand}(-c_{\text{max}}, c_{\text{max}}) \]  

Where, \( \text{rand} \) function is considered to generate the random numbers and \( c_{\text{max}} \) denotes the maximal change at current random perturbation.

Afterwards, it search the neighborhood nodes in \( K \) which has located near to \((\tilde{a}, \tilde{\beta})\), so on these nodes selects a novel set of CHs and denoted by \( \hat{L} \) (this can be infer as \( L_b = |\hat{L}| = |N| = |\hat{L}| \)). In present state is presented through \( \hat{L} \) set of CHs with \( f(\hat{L}) \) cost function and its neighborhood state is presented through \( \hat{L} \) set of CHs with \( f(\hat{L}) \) cost function.

At considered \( m \) iteration, the present state \( \hat{L} \) will change to \( \hat{L} \) neighborhood state and the probability can be computed as follows;

\[ N(\hat{L} \rightarrow L) = \exp[-f(\hat{L}) - f(L)/Q_m] \]  

The above given equation (11) is true when \( f(\hat{L}) \geq f(L) \), otherwise the probability of \( L \rightarrow \hat{L} \) will be one. Where, \( Q_m \) denotes for control factor parameter and given by;

\[ Q_m = R_{hp} \times e^{-\frac{m}{T0}} \]  

Where, \( R_{hp} \) is hyper parameter value that ranging from 100 to 10000. The computed value of \( Q_m \) help to decide the good clusters and function \( f(\cdot) \) signifies the cost function as;

\[ f(L) = \sum_{i=1}^{n} \min_{e_i \in E, l_i \in L} \vartheta(e_i, l_i) \]  

Where, distance in between the \( e_i \) node and \( l_i \) channel head is given by \( \vartheta(\cdot, \cdot) \). If \( L \) present state become \( \hat{L} \) neighborhood state then \( L = \hat{L} \).

In final level, \( S_{\text{itr}} \) stands for considered total iterations, where \( S_{\text{itr}} \) is greater or equal to \( m \) then algorithm is terminated otherwise it again go to neighborhood finding level. Once the determination of optimal CHs are done, the base station used to broadcast the message which consist of CH’s ID for every node present in a network. However, a particular node ID same to CH’s ID then it became a CH. Else the node selects its slot of TDMA for process of ‘data-transfer’ and node falls dead till the time has come to forward signal to its CH.

4. RESULTS AND DISCUSSION

In this section, the authors provided result and simulation analysis of our proposed OCEA with respect to other state-of-art technologies. The simulation is carried out in Matlab-2016b with system configuration of 12GB RAM, Intel i5 processor and latest Windows operating system. A WSN model with the sensor nodes is examined and it is assumed that our network model consists of 200 number of sensor nodes that uniformly scattered in the 100 x 100-meter monitoring area. The assumptions of the network model have been made where; the FNs and BS are nor doing the energy constrain, while the MNs are the energy-constrained. Moreover, each node knows its velocity, current position, and energy level and considered assumptions are reasonable for several real networks.

It is known that LEACH approach saves the energy of nodes in order to optimize efficiency of network in some aspects, but it lacks during complete network management. In addition, it is not much efficient to balance the energy consumption of network and to select the cluster heads sum. The OCEA algorithm periodically chooses the optimal set of the CHs according to an average node of hybrid energy and average speed of node.

To compare with our proposed approach, energy balance clustering routing algorithm (EB-CRA) is considered [32]. In [32], they randomly scattered 200 nodes in the 100m x 100m area, which is similar as per our simulation initialization. The EB-CRA is proposed for cluster head selection and BS communication, it considers the node residual energy and distance between the BS and CH. Therefore, the possibility of node energy is to be designated as the CH and the CH is used at the process of data transmission due to more distance from the BS. Moreover, the ACO algorithm [32] has perform better than the LEACH algorithm, so the authors consider this approach for CH energy consumption comparison.
Figure 2 shows the optimized cluster establishment and localization of CHs, where (a) represents at 1 poll, (b) represents at 10 poll and (c) represents at 500 poll. It is clearly noted that while increase in polling number or iteration or round, the number of clusters has increased but after several polling round it is not increasing. Which states that it has achieved is optimal cluster number. Increasing in number of polling, the localization of CHs is going to better and better.

CH nodes energy consumptions is very essential indicator to evaluate network efficiency and performance. Figure 3 shows energy consumption (Joules) of CH as per polling number. It shows that as per increment of polling, the energy consumption CHs in a network is decreasing. In case, polling round delay time is more compared to threshold value then the SN is considered to be dead node, Figure 4 represents the number of dead nodes as per polling number. Where it seems that till 160 polling number the dead nodes are zero and afterwards observed linear pattern between number of dead nodes (i.e., ranging from 0 to 60) and polling number (i.e., ranging from 160 to 500).

Figure 5 shows the average number of CHs as per polling number, this computed by taking mean CHs at each 100 polls. In Figure 5, from polls 1 to 100 and 101 to 200, the average number of CHs is same 37. While from 201 to 300, 301 to 400 and 401 to 500, the average number of CHs are 33, 31, and 27. It indicates the average of CHs count is decreasing as per increasing in number of polls.
Table 1 shows the energy comparison at different number of polls and same thing represented in a graphical format in Figure 6, where it considered different poll number to compare energy consumptions. At 100 poll number, our proposed methodology has got 11.35%, 64.7% and 45.8% less energy consumption compared to EB-CRA, LEACH and ACO techniques. Similarly, at 500 poll number, our proposed methodology has got 50.8%, 75.68% and 65.11% less energy consumption compared to EB-CRA, LEACH and ACO techniques. It indicates that difference in energy consumption has increased more with increment in polling number and our proposed methodology shows ability to manage less energy consumption at CHs.

| No. of Polls | ACO  | LEACH | EB-CRA | OCEA  |
|--------------|------|-------|--------|-------|
| 100          | 0.228| 0.3501| 0.1394 | 0.123576 |
| 200          | 0.2709| 0.3 | 0.1141 | 0.104796 |
| 300          | 0.2 | 0.2649 | 0.0791 | 0.087338 |
| 400          | 0.1723| 0.3248 | 0.1376 | 0.072483 |
| 500          | 0.175 | 0.2511 | 0.1243 | 0.061048 |

Figure 6. Graphical representation of energy comparison at different number of polls
5. CONCLUSION

In order to improve the lifetime and efficiency of the network, the computation of optimal cluster and localization of CHs is very important. Therefore in this study the authors proposed OCEA, which periodically chooses the optimal set of the CHs according to an average node of hybrid energy and average speed of node. In the result analysis section, shows the localization CHs in various polls, energy consumption with respect to polling round has also shown. It analyzed that CH nodes energy consumptions is decreasing as per polling round, which very essential indicator to evaluate network efficiency and performance. To compare with our proposed approach, EB-CRA, ACO and very popular LEACH protocol is considered in the validation part, where CH energy consumption is compared at various polls. Whereas, OCEA approach shows much effectiveness to balance the energy consumption of network and to select the cluster heads sum.

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