Abstract. Whereas set-valued tableaux are the combinatorial objects associated to stable Grothendieck polynomials, hook-valued tableaux are associated to stable canonical Grothendieck polynomials. In this paper, we define a novel uncrowding algorithm for hook-valued tableaux. The algorithm “uncrows” the entries in the arm of the hooks and yields a set-valued tableau and a column-flagged increasing tableau. We prove that our uncrowding algorithm intertwines with crystal operators. An alternative uncrowding algorithm that “uncrows” the entries in the leg instead of the arm of the hooks is also given. As an application of uncrowding, we obtain various expansions of the canonical Grothendieck polynomials.

1. Introduction

Set-valued tableaux play an important role in the $K$-theory of the Grassmannian. They form a generalization of semi-standard Young tableaux, where boxes may contain sets of integers rather than just integers [Buc02]. In particular, the stable symmetric Grothendieck polynomial indexed by the partition $\lambda$ is the generating function of set-valued tableaux

$$G_\lambda(x; \beta) = \sum_{T \in \text{SVT}(\lambda)} \beta^{|T|-|\lambda|_x}\text{weight}(T),$$

where $\text{SVT}(\lambda)$ is the set of set-valued tableaux of shape $\lambda$ and $\text{weight}(T)$ is the vector with $i$-th entry being the number of $i$ in $T$. Here $|T|$ is the number of entries in $T$ and $|\lambda|$ is the size of $\lambda$. Stable symmetric Grothendieck polynomials $G_\lambda$ can be viewed as a $K$-theory analogue of the Schur functions $s_\lambda$ (while the Grothendieck polynomial is an analog of the Schubert polynomial [LS83]). Buch [Buc02] also described the structure coefficients $c_{\lambda\mu}^{\nu}$, which is the coefficient of $G_\mu$ in the expansion of $G_\lambda G_\mu$ in terms of set-valued tableaux, generalizing the Littlewood–Richardson rule for Schur functions.

The Grassmannian $\text{Gr}(k, \mathbb{C}^n)$ of $k$-planes in $\mathbb{C}^n$ has a fundamental duality isomorphism

$$\text{Gr}(k, \mathbb{C}^n) \cong \text{Gr}(n - k, \mathbb{C}^n).$$

This implies that the structure constants have the symmetry $c_{\lambda\mu}^{\nu} = c_{\nu\lambda'}^{\mu'}$, where $\lambda'$ denotes the conjugate of the partition $\lambda$ (see for example [Ful97, Example 9.20]). Hence one expects a ring homomorphism on the completion of the ring of symmetric function defined on the basis of stable symmetric Grothendieck polynomials $\tau(G_\lambda) = G_{\lambda'}$. The standard involutive ring automorphism $\omega$ defined on the Schur basis by $\omega(s_\lambda) = s_{\lambda'}$ does not have this property [LP07]

$$\omega(G_\lambda) = J_\lambda \neq G_{\lambda'},$$

where $J_\lambda$ is the weak symmetric Grothendieck polynomial.

Yeliussizov [Yel17] introduced a new family of canonical stable Grothendieck polynomials $G_\lambda(x; \alpha, \beta)$ such that

$$\omega(G_\lambda(x; \alpha, \beta)) = G_{\lambda'}(x; \beta, \alpha).$$

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Combinatorially, the canonical stable Grothendieck polynomials can be expressed as generating functions of hook-valued tableaux. In a hook-valued tableau, each box contains a semistandard Young tableau of hook shape, which is weakly increasing in rows and strictly increasing in columns. More precisely

\[ G_\lambda(x; \alpha, \beta) = \sum_{T \in \text{HVT}(\lambda)} \alpha^{a(T)} \beta^{\ell(T)} x^{\text{weight}(T)}, \]

where HVT(\lambda) is the set of hook-valued tableaux of shape \lambda, a(T) is the sum of all arm lengths and \ell(T) is the sum of all leg lengths of the hook tableaux in T.

A hook-valued tableau T is a set-valued tableau when all hook tableaux entries are single columns or equivalently \( a(T) = 0 \). Hence \( G_\lambda(x; \alpha, \beta) \) specializes to \( G_\lambda(x; \beta) \) for \( \alpha = 0 \). Similarly, a hook-valued tableau T is a multiset-valued tableau when all hook tableaux entries are single rows or equivalently \( \ell(T) = 0 \). Hence \( G_\lambda(x; \alpha, \beta) \) specializes to \( J_\lambda(x; \alpha) \) for \( \beta = 0 \).

In this paper, we describe a novel uncrowding algorithm on hook-valued tableaux (see Definitions 3.2, 3.4 and 3.5). The uncrowding algorithm on set-valued tableaux was originally developed by Buch \[ Buc02, \text{Theorem 6.11} \] to give a bijective proof of Lenart’s Schur expansion of symmetric stable Grothendieck polynomials \[ Len00 \]. This uncrowding algorithm takes as input a set-valued tableau and produces a semistandard Young tableau (using the RSK bumping algorithm to uncrowd cells that contain more than one integer) and a flagged increasing tableau \[ Len00 \] (also known as an elegant filling \[ LP07, BM12, Pat16 \]), which serves as a recording tableau.

Chan and Pflueger \[ CP21 \] provide an expansion of stable Grothendieck polynomials indexed by skew partitions in terms of skew Schur functions. Their proof uses a generalization of the uncrowding algorithm of Lenart \[ Len00 \], Buch \[ Buc02 \], and Reiner, Tenner and Yong \[ RTY18 \] to skew shapes. Their analysis is motivated geometrically by identifying Euler characteristics of Brill–Noether varieties up to sign as counts of set-valued standard tableaux. The uncrowding algorithm was also used in the analysis of K-theoretic analogues of the Hopf algebras of symmetric functions, quasisymmetric functions, noncommutative symmetric functions, and of the Malvenuto–Reutenauer Hopf algebra of permutations \[ LP07, BM12, Pat16 \]. In \[ GZJ20 \], a vertex model for canonical Grothendieck polynomials and their duals was studied, which was used to derive Cauchy identities.

An important property of the uncrowding algorithm on set-valued tableaux is that it intertwines with crystal operators \[ MPS21 \] (see also \[ MPPS20 \]). The crystal structure on a combinatorial set is the combinatorial shadow of a (quantum) group representation (see for example \[ HK02, BS17 \]). A crystal structure on hook-valued tableaux was recently introduced by Hawkes and Scrimshaw \[ HS20 \]. Our novel uncrowding map on hook-valued tableaux yields a set-valued tableau and a recording tableau. We prove that it intertwines with crystal operators (see Proposition 3.12 and Theorem 3.14). This was stated as an open problem in \[ HS20 \].

The paper is organized as follows. In Section 2, we review the definition of semistandard hook-valued tableaux of \[ Yel17 \] and the crystal structure on them \[ HS20 \]. In Section 3, we define the new uncrowding map on hook-valued tableaux and prove that it intertwines with the crystal operators and other properties. We also give a variant of the uncrowding algorithm on hook-valued tableaux. In Section 4, we consider applications of the uncrowding algorithm, in particular expansions of the canonical Grothendieck polynomials using techniques developed in \[ BM12 \].

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### 2. Hook-valued tableaux

In Section 2.1, we define hook-valued tableaux \[ Yel17 \] and in Section 2.2 we review the crystal structure on hook-valued tableaux as introduced in \[ HS20 \].
2.1. **Hook-valued tableaux.** A semistandard Young tableau $U$ of hook shape is a tableau of the form

$$U = \begin{array}{c}
\ell_p \\
\vdots \\
\ell_1 \\
x & a_1 & \ldots & a_q
\end{array},$$

where the integer entries weakly increase from left to right and strictly increase from bottom to top. Note that we use French notation for Young diagrams and tableaux throughout the paper. In this case, $H(U) = x$ is called the **hook entry** of $U$, $L(U) = (\ell_1, \ell_2, \ldots, \ell_p)$ is the **leg** of $U$, and $A(U) = (a_1, a_2, \ldots, a_q)$ is the **arm** of $U$. Both the arm and the leg of $U$ are allowed to be empty. Additionally, the **extended leg** of $U$ is defined as $L^+(U) = (x, \ell_1, \ell_2, \ldots, \ell_p)$. We denote by $\max(U)$ (resp. $\min(U)$) the maximal (resp. minimal) entry in $U$.

**Definition 2.1.** [Yel17] Fix a partition $\lambda$. A **semistandard hook-valued tableau** (or **hook-valued tableau** for short) $T$ of shape $\lambda$ is a filling of the Young diagram for $\lambda$ with (nonempty) semistandard Young tableaux of hook shape such that:

(i) $\max(A) \leq \min(B)$ whenever the cell containing $A$ is in the same row, but left of the cell containing $B$;

(ii) $\max(A) < \min(C)$ whenever the cell containing $A$ is in the same column, but below the cell containing $C$.

The set of all hook-valued tableaux of shape $\lambda$ (respectively, with entries at most $m$) is denoted by $\text{HVT}(\lambda)$ (respectively, $\text{HVT}^m(\lambda)$).

Given a hook-valued tableau $T$, its **arm excess** is the total number of integers in the arms of all cells of $T$, while its **leg excess** is the total number of integers in the legs of all cells of $T$.

**Remark 2.2.** In the special case when a hook-valued tableau has arm excess 0, it is also called a **set-valued tableau**. Similarly, a **multiset-valued tableau** is a hook-valued tableau with leg excess 0. We use the notation $\text{SVT}(\lambda)$ (resp. $\text{SVT}^m(\lambda)$) and $\text{MVT}(\lambda)$ (resp. $\text{MVT}^m(\lambda)$) for the set of all set-valued tableaux of shape $\lambda$ (resp. with entries at most $m$) and the set of all multiset-valued tableaux of shape $\lambda$ (resp. with entries at most $m$), respectively.

2.2. **Crystal structure on hook-valued tableaux.** Hawkes and Scrimshaw [HS20] defined a crystal structure on hook-valued tableaux. We review their definition here.

**Definition 2.3** ([HS20], Definition 4.1). Let $C$ be a hook-valued tableau of column shape. The column reading word $R(C)$ is obtained by reading the extended leg in each cell from top to bottom, followed by reading all of the remaining entries, arranged in a weakly increasing order.

For a hook-valued tableau $T$, its column reading word is formed by concatenating the column reading words of all of its columns, read from left to right, that is,

$$R(T) = R(C_1)R(C_2)\ldots R(C_\ell),$$

where $\ell$ is the number of columns of $T$ and $C_i$ is the $i$th column of $T$.

**Example 2.4.** Let $T$ be the hook-valued tableau

$$T = \begin{array}{c|c|c}
4 & 33 & 5 \\
2 & 4 & \\
11 & 334 & 4445
\end{array}.$$

The column reading words for the columns of $T$ are respectively $432113$, $54334$ and $4445$, so that $R(T) = 432113543344445$. 
Definition 2.5. [HS20, Definition 4.3] Let \( T \in \text{HVT}^m(\lambda) \). For any \( 1 \leq i < m \), we employ the following pairing rules. Assign \( - \) to every \( i \) in \( R(T) \) and assign \( + \) to every \( i + 1 \) in \( R(T) \). Then, successively pair each \( + \) that is adjacent to and to the left of a \( - \), removing all paired signs until nothing can be paired.

The operator \( f_i \) acts on \( T \) according to the following rules in the given order. If there is no unpaired \( - \), then \( f_i \) annihilates \( T \). Otherwise, locate the cell \( c \) with entry the hook-valued tableau \( B = T(c) \) containing the \( i \) corresponding to the rightmost unpaired \( - \).

(M) If there is an \( i + 1 \) in the cell above \( c \) with entry \( B^{\uparrow} \), then \( f_i \) removes an \( i \) from \( A(B) \) and adds \( i + 1 \) to \( A(B^{\uparrow}) \).

(S) Otherwise, if there is a cell to the right of \( c \) with entry \( B^{\rightarrow} \), such that it contains an \( i \) in \( L^+(B^{\rightarrow}) \), then \( f_i \) removes the \( i \) from \( L^+(B^{\rightarrow}) \) and adds \( i + 1 \) to \( L(B) \).

(N) Else, \( f_i \) changes the \( i \) in \( B \) into an \( i + 1 \).

Similarly, the operator \( e_i \) acts on \( T \) according to the following rules in the given order. If there is no unpaired \( + \), then \( e_i \) annihilates \( T \). Otherwise, locate the cell \( c \) with entry the hook-valued tableau \( B = T(c) \) containing the entry \( i + 1 \) corresponding to the leftmost unpaired \( + \).

(M) If there is an \( i \) in the cell below \( c \) with entry \( B^{\downarrow} \), then \( e_i \) removes the \( i + 1 \) from \( A(B) \) and adds \( i \) to \( A(B^{\downarrow}) \).

(S) Otherwise, if there is a cell to the left of \( c \) with entry \( B^{\leftarrow} \), such that it contains an \( i + 1 \) in \( L(B^{\leftarrow}) \), then \( e_i \) removes the \( i + 1 \) from \( L(B^{\leftarrow}) \) and adds \( i \) to \( L^+(B) \).

(N) Else, \( e_i \) changes the \( i + 1 \) in \( B \) into an \( i \).

Based on the pairing procedure above, \( \varphi_i(T) \) is the number of unpaired \( - \), whereas \( \varepsilon_i(T) \) is the number of unpaired \( + \).

We remark that the definition of crystal operators on HVT specializes to the definition on SVT in [MPS21] or the one on MVT in [HS20] when the arm excess or leg excess of the tableaux is set to 0, respectively.

Example 2.6. Consider the following hook-valued tableau \( T \):

\[
T = \begin{bmatrix}
4 & 5 \\
34 & 4 \\
2 & 3 \\
11 & 233
\end{bmatrix}
\]

Then, \( e_3 \) annihilates \( T \), whereas

\[
e_1(T) = \begin{bmatrix}
4 & 5 \\
34 & 4 \\
3 & 2 \\
11 & 133
\end{bmatrix}, \quad f_1(T) = \begin{bmatrix}
4 & 5 \\
34 & 4 \\
2 & 3 \\
12 & 233
\end{bmatrix}, \quad f_3(T) = \begin{bmatrix}
4 & 5 \\
34 & 44 \\
2 & 3 \\
11 & 23
\end{bmatrix}
\]

For a given cell \((r, c)\) in row \( r \) and column \( c \) in a hook-valued tableau \( T \), let \( L_T(r, c) \) be the leg of \( T(r, c) \), let \( A_T(r, c) \) be arm of \( T(r, c) \), let \( H_T(r, c) \) be the hook entry of \( T(r, c) \), and let \( L_T^+(r, c) \) be the extended leg of \( T(r, c) \).

3. Uncrowding map on hook-valued tableaux

In Section 3.1, we first review the uncrowding map on set-valued tableaux. In Section 3.2, we give a new uncrowding map on hook-valued tableaux and prove some of its properties in Section 3.3. The relation to the uncrowding map on multiset-valued tableaux is given in Section 3.4. In Section 3.5, we give the inverse of the uncrowding map on hook-valued tableaux, called the crowding map. In Section 3.6, an alternative definition of the uncrowding map on hook-valued tableaux is provided.
3.1. Uncrowding map on set-valued tableaux. For set-valued tableaux, there exists an uncrowding operator, which maps a set-valued tableau to a pair of tableaux, one being a semistandard Young tableau and the other a flagged increasing tableau (see for example [Len00, Buc02, BM12, RTY18]). In this setting, the uncrowding operator intertwines with the crystal operators on set-valued tableaux and semistandard Young tableaux, respectively [MPS21].

Consider partitions $\lambda, \mu$ with $\lambda \subseteq \mu$ and $\lambda_1 = \mu_1$. A flagged increasing tableau (introduced in [Len00] and called (strict) elegant fillings by various authors [LP07, BM12, Pat16]) is a row and column strict filling of the skew shape $\mu/\lambda$ such that the positive integer entries in the $i$-th row of the tableau are at most $i - 1$ for all $1 \leq i \leq \ell(\mu)$, where $\ell(\mu)$ is the length of partition $\mu$. In particular, the bottom row is empty. The set of all flagged increasing tableaux is denoted by $\mathcal{F}$. The set of all flagged increasing tableaux of shape $\mu/\lambda$ with $\lambda_1 = \mu_1$ is denoted by $\mathcal{F}(\mu/\lambda)$.

We now review the uncrowding operation on set-valued tableaux. We call a cell in a set-valued tableau a multicell if it contains more than one letter.

Definition 3.1. Define the uncrowding operation on $T \in \text{SVT}(\lambda)$ as follows. First identify the topmost row $r$ in $T$ with a multicell. Let $x$ be the largest letter in row $r$ that lies in a multicell; remove $x$ from the cell and perform RSK row bumping with $x$ into the rows above. The resulting tableau, whose shape differs from $\lambda$ by the addition of one cell, is the output of this operation.

The uncrowding map on set-valued tableaux

$$ U_{\text{SVT}} : \text{SVT}(\lambda) \rightarrow \bigsqcup_{\mu \geq \lambda} \text{SSYT}(\mu) \times \mathcal{F}(\mu/\lambda) $$

is defined as follows. Let $T \in \text{SVT}(\lambda)$ with leg excess $\ell$.

1. Initialize $P_0 = T$ and $Q_0 = F_0$, where $F_0$ is the unique flagged increasing tableau of shape $\lambda/\lambda$.
2. For each $1 \leq i \leq \ell$, $P_i$ is obtained from $P_{i-1}$ by applying the uncrowding operation. Let $C$ be the cell in $\text{shape}(P_i)/\text{shape}(P_{i-1})$. If $C$ is in row $r'$, then $F_i$ is obtained from $F_{i-1}$ by adding cell $C$ with entry $r' - r$.
3. Set $U_{\text{SVT}}(T) = (P, F) := (P_\ell, F_\ell)$.

It was proved in [Buc02, Section 6] that $U_{\text{SVT}}$ in (3.1) is a bijection. Monical, Pechenik and Scrimshaw [MPS21] proved that $U_{\text{SVT}}$ intertwines with the crystal operators on set-valued tableaux (see also [MPPS21]). A similar uncrowding algorithm for multiset-valued tableaux was given in [HS20, Section 3.2].

3.2. Uncrowding map on hook-valued tableaux. In [HS20], the authors ask for an uncrowding map for hook-valued tableaux which intertwines with the crystal operators. Here we provide such an uncrowding map by uncrowding the arm excess in a hook-valued tableau to obtain a set-valued tableau. An alternative obtained by uncrowding the leg excess first is given in Section 3.4.

Definition 3.2. The uncrowding bumping $\mathcal{V}_{\partial} : \text{HVT} \rightarrow \text{HVT}$ is defined by the following algorithm:

1. Initialize $T$ as the input.
2. If the arm excess of $T$ equals zero, return $T$.
3. Else, find the rightmost column that contains a cell with nonzero arm excess. Within this column, find the cell with the largest value in its arm. (In French notation this is the topmost cell with nonzero arm excess in the specified column.) Denote the row index and column index of this cell by $r$ and $c$, respectively. Denote the cell as $(r, c)$, its rightmost arm entry by $a$, and its largest leg entry by $\ell$.
4. Look at the column to the right of $(r, c)$ (i.e. column $c + 1$) and find the smallest number that is greater than or equal to $a$.
   - If no such number exists, attach an empty cell to the top of column $c + 1$ and label the cell as $(\tilde{r}, c + 1)$, where $\tilde{r}$ is its row index. Let $k$ be the empty character.
If such a number exists, label the value as \( k \) and the cell containing \( k \) as \((\tilde{r}, c+1)\) where \( \tilde{r} \) is the cell’s row index.

We now break into cases:

(a) If \( \tilde{r} \neq r \), then remove \( a \) from \( A_T(r, c) \), replace \( k \) with \( a \), and attach \( k \) to the arm of \( A_T(\tilde{r}, c+1) \).

(b) If \( \tilde{r} = r \) then remove \((a, \ell) \cap L_T(r, c)\) from \( L_T(r, c) \) where \((a, \ell) = \{a + 1, a + 2, \ldots, \ell\}\), remove \( a \) from \( A_T(r, c) \), insert \((a, \ell) \cap L_T(r, c)\) into \( L_T(\tilde{r}, c+1) \), replace the hook entry of \((\tilde{r}, c+1)\) with \( a \), and attach \( k \) to \( A_T(\tilde{r}, c+1) \).

(5) Output the resulting tableau.

See Figures 1 and 2 for illustration.

\[
\begin{array}{c|c|c}
\hline
\ell & \ast & - \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
\end{array} 
\begin{array}{c|c|c}
\hline
- & - \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
\end{array} \rightarrow 
\begin{array}{c|c|c}
\hline
\ell & \ast & - \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
\end{array}
\]

Figure 1. When \( \tilde{r} \neq r \). Left: \((\tilde{r}, c+1)\) is a new cell; Right: \((\tilde{r}, c+1)\) is an existing cell.

\[
\begin{array}{c|c|c}
\hline
- & - a \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
\hline
\end{array} \rightarrow 
\begin{array}{c|c|c}
\hline
- & - \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
\hline
\end{array} \rightarrow 
\begin{array}{c|c|c}
\hline
- & - k \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
- & - \vspace{0.5cm} \\
\hline
\end{array}
\]

Figure 2. When \( \tilde{r} = r \). Left: \((r, c+1)\) is a new cell; Right: \((r, c+1)\) is an existing cell.

**Lemma 3.3.** The map \( \mathcal{V}_b \) is well-defined. More precisely, for \( T \in HVT \) we have \( \mathcal{V}_b(T) \in HVT \).

**Proof.** It suffices to check that \( \mathcal{V}_b \) preserves the semistandardness condition of both the entire hook-valued tableau and the filling within each cell. We break into two cases depending on whether Step (4)a or (4)b in Definition 3.2 is applied.

**Case 1:** Assume Step (4)a is applied. To verify semistandardness within each cell, it suffices to check cells \((r, c)\) and \((\tilde{r}, c+1)\). The semistandardness within cell \((r, c)\) is clearly preserved as the only change to the hook-shaped tableau in cell \((r, c)\) is that an entry was removed from \( A_T(r, c) \). We now check the semistandardness condition within cell \((\tilde{r}, c+1)\). We have that \( \mathcal{V}_b \) either created the cell \((\tilde{r}, c+1)\) and inserted the number \( a \) in it or \( \mathcal{V}_b \) replaced \( k \) with \( a \) and appended \( k \) to the arm of cell \((\tilde{r}, c+1)\). In both cases, the tableau in cell \((\tilde{r}, c+1)\) is a semistandard hook-shaped tableau. In the second case this is true since \( k \) is weakly greater than \( H_T(\tilde{r}, c+1) \) and \( k \) is the smallest number weakly greater than \( a \) in column \( c+1 \).

We now check the semistandardness of the entire tableau. Note that it suffices to check the semistandardness in row \( \tilde{r} \) and column \( c+1 \). Since \( \tilde{r} < r \), the semistandardness in row \( \tilde{r} \) is preserved as \( a \) is larger than every number in \((\tilde{r}, c)\) and \( k \) remains in the same cell. Also, the semistandardness in column \( c+1 \) is preserved as \( k \) is chosen to be the smallest number in column \( c+1 \) that is weakly greater than \( a \).

**Case 2:** Assume Step (4)b is applied. The semistandardness within cell \((r, c)\) is clearly preserved as the only change to \((r, c)\) is that entries from \( L_T(r, c) \) and \( A_T(r, c) \) are removed. We now check the semistandardness condition within cell \((r, c+1)\). If \((a, \ell) \cap L_T(r, c) = \emptyset \),
then $a$ is weakly larger than all elements of $(r, c)$. In this case, the semistandardness within cell $(r, c + 1)$ follows from the argument in Case 1. If $(a, ℓ] \cap L_T(r, c) \neq \emptyset$, then $a$ is not weakly larger than all elements of $(r, c)$. After applying $V_b$ the semistandardness condition in the leg of $(r, c + 1)$ will still hold as $a < x < z$ for all $x \in (a, ℓ] \cap L_T(r, c)$, where $z$ is the smallest value in $L_T(r, c + 1)$. Similarly, the semistandardness condition in the arm of $(r, c + 1)$ holds as $a < k$ or $k$ is the empty character. Thus, the semistandardness condition in each cell is preserved. The semistandardness of row $r$ is preserved as all numbers strictly greater than $a$ in $(r, c)$ are moved to $(r, c + 1)$ along with $a$. The semistandardness condition within column $c + 1$ is preserved as every number in $(r + 1, c + 1)$ is strictly greater than $ℓ$ and every number in $(r - 1, c + 1)$ is strictly less than $a$.

\[\square\]

**Definition 3.4.** The uncrowding insertion $V : \text{HVT} \to \text{HVT}$ is defined as $V(T) = V^d_b(T)$, where the integer $d \geq 1$ is minimal such that $\text{shape}(V^d_b(T)) / \text{shape}(V^{d-1}_b(T)) \neq \emptyset$ or $V^d_b(T) = V^{d-1}_b(T)$.

A column-flagged increasing tableau is a tableau whose transpose is a flagged increasing tableau. Let $F^c$ denote the set of all column-flagged increasing tableaux. Let $F^c(\mu/\lambda)$ denote the set of all column-flagged increasing tableaux of shape $\mu/\lambda$.

**Definition 3.5.** Let $T \in \text{HVT}(\lambda)$ with arm excess $\alpha$. The uncrowding map

$$U : \text{HVT}(\lambda) \to \bigsqcup_{\mu \supseteq \lambda} \text{SVT}(\mu) \times F^c(\mu/\lambda)$$

is defined by the following algorithm:

1. Let $P_0 = T$ and let $Q_0$ be the column-flagged increasing tableau of shape $\lambda/\lambda$.

2. For $1 \leq i \leq \alpha$, let $P_{i+1} = V(P_i)$. Let $c$ be the index of the rightmost column of $P_i$ containing a cell with nonzero arm excess and let $\tilde{c}$ be the column index of the cell $\text{shape}(P_{i+1}) / \text{shape}(P_i)$. Then $Q_{i+1}$ is obtained from $Q_i$ by appending the cell $\text{shape}(P_{i+1}) / \text{shape}(P_i)$ to $Q_i$ and filling this cell with $\tilde{c} - c$.

Define $U(T) = (P(T), Q(T)) := (P_\alpha, Q_\alpha)$.

**Example 3.6.** Let $T$ be the hook-valued tableau

\[
\begin{array}{cccc}
8 & 67 \\
5 & 4 \\
233 & 66 \\
1 & 2 & 7 \\
& 11 & 5
\end{array}
\]

Then, we obtain the following sequence of tableaux $V^i_b(T)$ for $0 \leq i \leq 2 = d$ when computing the first uncrowding insertion:

\[
\begin{array}{c}
\begin{array}{c}
8 & 67 \\
5 & 4 \\
233 & 66 \\
1 & 2 & 7
\end{array} \\
\rightarrow \\
\begin{array}{c}
8 & 67 \\
5 & 4 \\
233 & 6 \\
1 & 2 & 6
\end{array} \\
\rightarrow \\
\begin{array}{c}
8 & 67 \\
5 & 4 \\
233 & 6 \\
1 & 2 & 6
\end{array} = V(T).
\end{array}
\]
Continuing with the remaining uncrowding insertions, we obtain the following sequences of tableaux for the uncrowding map:

\[
\begin{array}{c|c|c}
8 & 67 \\
5 & 4 \\
233 & 66 \\
1 & 2 & 7 \\
\end{array}
\rightarrow
\begin{array}{c|c|c}
5 & 67 \\
4 & 3 \\
233 & 6 \\
1 & 2 & 6 \\
\end{array}
\rightarrow
\begin{array}{c|c|c}
8 & 67 \\
5 & 4 \\
233 & 6 \\
1 & 2 & 6 \\
\end{array}
\rightarrow
\begin{array}{c|c|c}
8 & 67 \\
5 & 4 \\
233 & 6 \\
1 & 2 & 6 \\
\end{array}
= P(T),
\]

\[
\begin{array}{c|c|c}
6 & 7 \\
5 & 4 \\
233 & 6 \\
1 & 1 & 2 & 6 \\
\end{array}
\rightarrow
\begin{array}{c|c|c}
5 & 7 \\
4 & 2 \\
233 & 6 \\
1 & 2 & 6 \\
\end{array}
\rightarrow
\begin{array}{c|c|c}
5 & 7 \\
4 & 2 \\
233 & 6 \\
1 & 2 & 6 \\
\end{array}
\rightarrow
\begin{array}{c|c|c}
5 & 7 \\
4 & 2 \\
233 & 6 \\
1 & 1 & 1 & 5 & 6 & 7 \\
\end{array}
= Q(T).
\]

Corollary 3.7. Let \( T \in \text{HVT} \). Then \( P(T) \) is a set-valued tableau.

Proof. By Lemma 3.3 and Definition 3.4, we have that \( \mathcal{V}(T) \) is a hook-valued tableau. Note that if the arm excess of \( T \) is nonzero, then the arm excess of \( \mathcal{V}(T) \) is one less than that of \( T \). Since \( P(T) = \mathcal{V}^\alpha(T) \), where \( \alpha \) is the arm excess of \( T \), we have that the arm excess of \( P(T) \) is zero. Thus, \( P(T) \) is a set-valued tableau. \( \square \)

Definition 3.8. Let \( T \in \text{HVT} \) and let \( d \) be minimal such that \( \mathcal{V}(T) = \mathcal{V}_b^d(T) \). The insertion path \( p \) of \( T \rightarrow \mathcal{V}(T) \) is defined as follows:

- If \( d = 0 \), set \( p = \emptyset \).
- Otherwise, let \( (r_0, c_0) \) be the rightmost and topmost cell of \( T \) containing a cell with nonzero arm excess. For all \( 1 \leq j \leq d \), let \( c_j = c_0 + j \) and let \( r_j = \tilde{r} \) be \( \tilde{r} \) in Definition 3.2 when \( \mathcal{V}_b \) is applied to \( \mathcal{V}_b^{j-1}(T) \). Set \( p = ((r_0, c_0), (r_1, c_1), \ldots, (r_d, c_d)) \).

Lemma 3.9. Let \( T \in \text{HVT} \). Then \( Q(T) \) is a column-flagged increasing tableau.

Proof. By construction, the positive integer entries in column \( i \) of \( Q(T) \) are at most \( i - 1 \). Let \( m \) be the smallest nonnegative integer such that \( \mathcal{V}^m(T) = P(T) \). Let \( p^i = ((r_0^i, c_0^i), (r_1^i, c_1^i), \ldots, (r_d^i, c_d^i)) \) for \( 0 \leq i < m \) be the insertion path of \( \mathcal{V}^i(T) \rightarrow \mathcal{V}^{i+1}(T) \). Since \( c_{i+1}^0 \leq c_0^i \) for all \( 0 \leq i < m \), the entries in each row of \( Q(T) \) are strictly increasing. To check that the entries in each column of \( Q(T) \) are strictly increasing, it suffices to show that if \( c_{i+1}^0 = c_0^i \), then \( p_{i+1} \) lies weakly below \( p_i \). In other words, it suffices to check that \( c_0^{i+1} = c_0^i \) implies that \( r_{j+1}^{i+1} \leq r_j^i \) for all \( 0 \leq j \leq d_i \). We prove this by induction on \( j \). Note that \( r_{i+1}^0 \leq r_0^i \) by the definition of \( \mathcal{U} \). Assume by induction that \( r_{j+1}^{i+1} \leq r_j^i \). This implies that the \( a \) when applying \( \mathcal{V}_b \) to \( \mathcal{V}_b^i(\mathcal{V}^j(T)) \) is weakly smaller than the \( a \) when applying \( \mathcal{V}_b \) to \( \mathcal{V}_b^j(\mathcal{V}^{j-1}(T)) \). Thus, we must have \( r_{j+1}^{i+1} \leq r_{j+1}^i \). \( \square \)
3.3. Properties of the uncrowding map. Let \( T \) be a hook-valued tableau. Define \( R_i(T) \) as the induced subword of \( R(T) \) consisting only of the letters \( i \) and \( i + 1 \). In the next lemma, we use the same notation as in Definition 3.2. Furthermore, two words are Knuth equivalent if one can be transformed to the other by a sequence of Knuth equivalences on three consecutive letters

\[
xyz \equiv zyx \quad \text{for } x \leq y < z, \quad yxz \equiv yzx \quad \text{for } x < y \leq z.
\]

Lemma 3.10. For \( T \in \text{HVT} \), \( R_i(T) = R_i(\mathcal{V}_b(T)) \) unless \( T \) satisfies one of the following three conditions:

(a) \( a = i \) or \( a = i + 1 \) and column \( c + 1 \) contains both \( i \) and \( i + 1 \),

(b) \( \tilde{r} = r, i \in (a, \ell] \cap L_T(r, c), k = i, \) and column \( c + 1 \) contains \( i + 1 \),

(c) \( \tilde{r} = r, a = i, i + 1 \in (a, \ell] \cap L_T(r, c), \) and \( (r, c) \) contains another \( i \) besides \( a \).

Moreover, \( R_i(T) \) is Knuth equivalent to \( R_i(\mathcal{V}_b(T)) \).

Proof. Let \( R_i(T) = r_1 r_2 \ldots r_m \). We break into cases based on the value of \( a \).

Case 1: Assume \( a \neq i, i + 1 \).

Assume Step (4)a is applied by \( \mathcal{V}_b \). If \( k \neq i, i + 1 \), then \( R_i(T) = R_i(\mathcal{V}_b(T)) \) as the position of all letters \( i \) and \( i + 1 \) remains the same. Let \( k = i \). We have that \( k \) is the only \( i \) in column \( c + 1 \). Hence, when \( k \) gets bumped from \( L_T(\tilde{r}, c + 1) \) and appended to \( A_T(\tilde{r}, c + 1) \), the relative position of \( k \) to the other letters \( i \) and \( i + 1 \) in \( R_i(T) \) does not change. Thus, \( R_i(T) = R_i(\mathcal{V}_b(T)) \). Let \( k = i + 1 \). Note that column \( c + 1 \) cannot have a cell containing an \( i \) as \( k \) is the smallest number weakly greater than \( a \). Hence, moving \( k \) from \( L_T(\tilde{r}, c + 1) \) to \( A_T(\tilde{r}, c + 1) \) will not change \( R_i(T) \). Therefore, we once again have that \( R_i(T) = R_i(\mathcal{V}_b(T)) \).

Assume Step (4)b is applied by \( \mathcal{V}_b \). Consider the subcase when \( (a, \ell] \cap L_T(r, c) = \emptyset \). By a similar argument to the previous paragraph, we have that \( R_i(T) = R_i(\mathcal{V}_b(T)) \). Next, consider the subcase when \( i + 1 \in (a, \ell] \cap L_T(r, c) \). This implies that \( a < i \) and the only time \( i + 1 \) occurs in column \( c \) is in \( L_T(r, c) \). Note that if an \( i \) exists in column \( c \), it must be contained in \( L_T(r, c) \). We also have that \( k \geq i + 1 \) or \( k \) is the empty character and no cell in column \( c + 1 \) contains an \( i \). Thus, removing \( (a, \ell] \cap L_T(r, c) \) from \( L_T(r, c) \), replacing \( k \) with \( (a, \ell] \cap L_T(r, c) \) in \( L_T(r, c + 1) \), and appending \( k \) to \( A_T(r, c + 1) \) does not change \( R_i(T) \). Therefore \( R_i(T) = R_i(\mathcal{V}_b(T)) \). Let \( i \in (a, \ell] \cap L_T(r, c) \) and \( i + 1 \notin (a, \ell] \cap L_T(r, c) \). Note that the only place \( i + 1 \) can occur in column \( c \) is as \( H_T(r + 1, c) \) and the only place \( i \) can occur in column \( c \) is in \( L_T(r, c) \). This implies that removing \( (a, \ell] \cap L_T(r, c) \) from \( L_T(r, c) \), replacing \( k \) with \( (a, \ell] \cap L_T(r, c) \) in \( L_T(r, c + 1) \) and appending \( k \) to \( A_T(r, c + 1) \) will not change \( R_i(T) \) unless both \( i + 1 \) and \( i \) show up in column \( c + 1 \). This can only occur when \( k = i \) which implies that \( R_i(T) = r_1 \ldots i \ i + 1 \ k \ldots r_m \) and \( R_i(\mathcal{V}_b(T)) = r_1 \ldots i + 1 \ i \ k \ldots r_m \). We see that \( R_i(T) \) and \( R_i(\mathcal{V}_b(T)) \) only differ by a Knuth relation implying they are Knuth equivalent. Assume that \( i, i + 1 \notin (a, \ell] \cap L_T(r, c) \) \( \neq \emptyset \). If \( a > i + 1 \) the positions of all letters \( i \) and \( i + 1 \) remain the same after \( \mathcal{V}_b \) is applied. If \( a < i \), then the positions of all letters \( i \) and \( i + 1 \) also remain the same unless \( k = i \) or \( k = i + 1 \). In both of these special subcases, it can be checked that still \( R_i(T) = R_i(\mathcal{V}_b(T)) \).

Case 2: Assume \( a = i \).

Assume Step (4)a is applied by \( \mathcal{V}_b \). If column \( c + 1 \) does not contain both \( i \) and \( i + 1 \), then we have \( R_i(T) = R_i(\mathcal{V}_b(T)) \). However, if both an \( i \) and an \( i + 1 \) are in column \( c + 1 \), then \( R_i(T) = r_1 \ldots i + 1 \ i \ldots r_m \) and \( R_i(\mathcal{V}_b(T)) = r_1 \ldots i + 1 \ i \ldots r_m \) which are Knuth equivalent.

Assume Step (4)b is applied by \( \mathcal{V}_b \). Consider the subcase when \( (a, \ell] \cap L_T(r, c) = \emptyset \). By a similar argument to the previous paragraph, we have that \( R_i(T) = R_i(\mathcal{V}_b(T)) \) unless both an \( i \) and an \( i + 1 \) are in column \( c + 1 \) in which case \( R_i(T) \) and \( R_i(\mathcal{V}_b(T)) \) are only Knuth equivalent. Consider the subcase given by \( i + 1 \in (a, \ell] \cap L_T(r, c) \). Note that no cell in column \( c + 1 \) can contain an \( i \), the only cell that could contain an \( i + 1 \) in column \( c + 1 \) is \( (r, c + 1) \), and the only cell containing letters \( i \) or \( i + 1 \) in column \( c \) is \( (r, c) \). This implies that it suffices to look at the changes to \( (r, c) \) and \( (r, c + 1) \). We see that \( R_i(T) = r_1 \ldots i + 1 \ \ldots i a \ldots r_m \) and \( R_i(\mathcal{V}_b(T)) = r_1 \ldots i \ldots i + 1 \ a \ldots r_m \).
where $\gamma \geq 1$ is the number of letters $i$ in cell $(r, c)$ including $a$. We see that $R_i(T)$ and $R_i(V_b(T))$ are Knuth equivalent. Consider the subcase when $i + 1 \not\in (a, \ell] \cap L_T(r, c) \neq \emptyset$. We have that both $i$ and $i + 1$ cannot be in a cell in column $c + 1$ and an $i + 1$ cannot be in column $c$. Thus applying $V_b$ does not change $R_i(T)$ giving us that $R_i(T) = R_i(V_b(T))$.

**Case 3:** Assume $a = i + 1$.

Assume Step (4)a is applied by $V_b$. If column $c + 1$ does not contain both $i$ and $i + 1$, then we have that $R_i(T) = R_i(V_b(T))$. However, if both $i$ and $i + 1$ occur in column $c + 1$, then $R_i(T) = r_1 \ldots i + 1 i + 1 i \ldots r_m$ and $R_i(V_b(T)) = r_1 \ldots i + 1 i + 1 i \ldots r_m$ which are Knuth equivalent.

Assume Step (4)b is applied by $V_b$. If $(a, \ell] \cap L_T(r, c) = \emptyset$, then $R_i(T) = R_i(V_b(T))$ unless both $i$ and $i + 1$ occur in column $c + 1$. In this exceptional case, we have that $R_i(T)$ and $R_i(V_b(T))$ are only Knuth equivalent by a similar argument to the previous paragraph. If $(a, \ell] \cap L_T(r, c) \neq \emptyset$, then $k > i + 1$ or $k$ is the empty character and no cell in column $c + 1$ contains an $i + 1$. Thus applying $V_b$ does not change $R_i(T)$ giving us that $R_i(T) = R_i(V_b(T))$. $\square$

**Remark 3.11.** In general, the full reading words are not Knuth equivalent under the uncrowding map. For example, take the following hook-valued tableau $T$, which uncrowds to a set-valued tableau $S$:

\[
T = \begin{array}{c|c|c|c}
4 & 3 & | & 5 \\
\hline
2 & 12 & 4 \\
\end{array} \rightarrow \begin{array}{c|c|c|c|c}
2 & 4 & 5 & | & 1 \\
\hline
4 & 3 & 5 & | & 4 \\
\end{array} = S.
\]

The reading word changed from 4321254 to 2143254, which are not Knuth equivalent.

**Proposition 3.12.** Let $T \in \text{HVT}$.

(1) If $f_i(T) = 0$, then $f_i(P(T)) = 0$.

(2) If $e_i(T) = 0$, then $e_i(P(T)) = 0$.

**Proof.** Since $P(T) = V_b^s(T)$ for some $s \in \mathbb{N}$ and Knuth equivalence is transitive, we have that $R_i(T)$ is Knuth equivalent to $R_i(P(T))$ by the previous lemma. As $f_i(T) = 0$, we have that every $i$ in $R_i(T)$ is $i$-paired with an $i + 1$ to its left. This property is preserved under Knuth equivalence giving us that $f_i(P(T)) = 0$. The same reasoning implies (2). $\square$

**Lemma 3.13.** Let $T \in \text{HVT}$.

(1) If $f_i(T) \neq 0$, then $f_i(V_b(T)) = V_b(f_i(T)) \neq 0$.

(2) If $e_i(T) \neq 0$, then $e_i(V_b(T)) = V_b(e_i(T)) \neq 0$.

**Proof.** We are going to prove (1). Part (2) follows since $e_i$ and $f_i$ are partial inverses.

Let $a, \ell, k, r, c$, and $\hat{r}$ be defined as in Definition 3.2 when $V_b$ is applied to $T$. Similarly, define $a', \ell', k', r', c'$, and $\hat{r}'$ for when $V_b$ is applied to $f_i(T)$. Let $R_i(T) = r_1 r_2 \ldots r_m$ and $R_i(V_b(T)) = r_1' r_2' \ldots r_m'$ be the corresponding reading words. Let $(\hat{r}, \hat{c})$ denote the cell containing the rightmost unpaired $i$ in $T$, where $\hat{r}$ and $\hat{c}$ are its row and column index respectively. We break into cases based on the position of $(\hat{r}, \hat{c})$ to $(r, c)$.

**Case 1:** Assume $(\hat{r}, \hat{c}) = (r, c)$. We break into subcases based on how $f_i$ acts on $T$.

- Assume that $(\hat{r} + 1, c)$ contains an $i + 1$.
  
  As every entry in $(r, c)$ must be strictly smaller than the values in $(r + 1, c)$ and $(r, c)$ must contain an $i$, we have that $\ell = i$ or $a = i$. If $\ell = i$, then $\ell$ is $i$-paired with the $i + 1$ in $(r + 1, c)$. Hence $a$ is always equal to $i$ and $a$ must correspond to the rightmost unpaired $i$ of $T$. Thus, $f_i$ acts on $T$ by removing $a$ from $(r, c)$ and appending an $i + 1$ to $A_T(r + 1, c)$. Note that $(a, \ell] \cap L_T(r, c) = \emptyset$ implying $V_b$ acts on $T$ by removing $a$ from $A_T(r, c)$, replacing $k$ in $(\hat{r}, c + 1)$ with $a$, and appending $k$ to $A_T(\hat{r}, c + 1)$ where $\hat{r} \leq r$. We break into subcases based upon where the values of $i$ and $i + 1$ are in column $c + 1$ utilizing the fact that column $c + 1$ cannot contain an $i$ without an $i + 1$ (since the arm
excess of cell \((r+1, c)\) is zero and cell \((r, c)\) contains the rightmost unpaired \(i\).

Assume that column \(c+1\) does not contain an \(i\). Since \(a\) corresponds to the rightmost unpaired \(i\) in \(T\) and column \(c+1\) does not contain an \(i\), we have that the rightmost unpaired \(i\) in \(V_h(T)\) is precisely \(a\) in the cell \((\tilde{r}, c+1)\). Note that \((\tilde{r}+1, c+1)\) does not contain an \(i+1\) in \(V_h(T)\) as \(k \geq i+1\) or \(k\) is the empty character. Similarly, we have that \((\tilde{r}, c+2)\) does not contain an \(i\). Thus, \(f_i\) acts on \(V_h(T)\) by changing \(a\) to an \(i+1\) in \((\tilde{r}, c+1)\). We now consider \(V_h(f_i(T))\). When applying \(V_h\) to \(f_i(T)\), \(a'\) is precisely the \(i+1\) appended to \(A_T(r+1, c)\) and \(k'\) is the same as \(k\). Since \(r' = \tilde{r} < r+1\), we have that \(V_h\) acts on \(f_i(T)\) by removing \(i+1\) from \(A_{f_i(T)}(r+1, c)\), replacing \(k\) with an \(i+1\) in \((\tilde{r}, c+1)\), and appending \(k\) to \(A_{f_i(T)}(\tilde{r}, c+1)\). We see that \(f_i(V_h(T)) = V_h(f_i(T))\).

Assume that column \(c+1\) contains both an \(i\) and an \(i+1\) in the same cell. Note that this implies that \(k = i\). Since \(a\) is the rightmost unpaired \(i\) in \(T\) and the only cell in column \(c+1\) that contained an \(i+1\) or an \(i\) is \((\tilde{r}, c+1)\), we have that the rightmost unpaired \(i\) in \(V_h(T)\) is the \(i\) appended to \(A_T(\tilde{r}, c+1)\). Since \((\tilde{r}, c+1)\) contains an \(i+1\), we have that \((\tilde{r}+1, c+1)\) cannot contain an \(i+1\) and \((\tilde{r}, c+2)\) cannot contain an \(i\). Thus, \(f_i\) acts on \(V_h(T)\) by changing the \(i\) in \(A_{V_h(T)}(\tilde{r}, c+1)\) to an \(i+1\). We now consider \(V_h(f_i(T))\). When applying \(V_h\) to \(f_i(T)\), \(a'\) is precisely the \(i+1\) appended to \(A_T(r+1, c)\) and \(k'\) is the \(i+1\) in \((\tilde{r}, c+1)\). Since \(r' = r < r+1\), we have that \(V_h\) acts on \(f_i(T)\) by removing \(i+1\) from \(A_{f_i(T)}(r+1, c)\), replacing \(i+1\) in \((\tilde{r}, c+1)\) with the \(i+1\) from \(A_{f_i(T)}(r+1, c)\), and appending an \(i+1\) to \(A_{f_i(T)}(\tilde{r}, c+1)\). We see that \(f_i(V_h(T)) = V_h(f_i(T))\).

Assume that column \(c+1\) contains both an \(i\) and an \(i+1\) in different cells. Note that this implies that \(k = i\). Since \(a\) corresponds to the rightmost unpaired \(i\) in \(R_i(T)\) and the only \(i+1\) and \(i\) in column \(c+1\) are in cells \((\tilde{r}+1, c+1)\) and \((\tilde{r}, c+1)\) respectively, we have that the rightmost unpaired \(i\) in \(R_i(V_h(T))\) corresponds to the \(i\) appended to \(A_T(\tilde{r}, c+1)\). By assumption, we have that \((\tilde{r}+1, c+1)\) contains an \(i+1\). Thus, \(f_i\) acts on \(V_h(T)\) by removing the \(i\) from \(A_{V_h(T)}(\tilde{r}, c+1)\) and appending an \(i+1\) to \(A_{V_h(T)}(\tilde{r}+1, c+1)\). We now consider \(V_h(f_i(T))\). When applying \(V_h\) to \(f_i(T)\), \(a'\) is precisely the \(i+1\) appended to \(A_T(r+1, c)\) and \(k'\) is the \(i+1\) in cell \((\tilde{r}+1, c+1)\). If \(r' = r+1\), then \(i+1\) is weakly larger than every value in \((r+1, c)\). Thus, either \((a', \ell') \cap L_{f_i(T)}(r+1, c) = \emptyset\) or \(r' < r+1\). This implies that \(V_h\) acts on \(f_i(T)\) by removing \(i+1\) from \(A_{f_i(T)}(r+1, c)\), replacing the \(i+1\) in \(H_{f_i(T)}(\tilde{r}+1, c+1)\) with the \(i+1\) removed from \(A_{f_i(T)}(r+1, c)\), and appending an \(i+1\) to \(A_{f_i(T)}(\tilde{r}+1, c+1)\). We see that \(f_i(V_h(T)) = V_h(f_i(T))\).

- Assume that \((r+1, c)\) does not contain an \(i+1\) and \((r, c+1)\) contains an \(i\).
  Under these assumptions, we have that no cell in column \(c\) can contain an \(i+1\). This implies that column \(c+1\) must contain an \(i+1\). The cell \((r+1, c+1)\) cannot have an \(i+1\) as this would force \((r+1, c)\) to also have an \(i+1\). Thus, \((r, c+1)\) must contain an \(i+1\) in its leg. By our assumption we have that \(f_i\) acts on \(T\) by removing the \(i\) from \((r, c+1)\) and appending an \(i+1\) to \(L_T(r, c)\). We break into subcases according to where the rightmost unpaired \(i\) sits inside the cell \((r, c)\). If the rightmost unpaired \(i\) is in \(H_T(r, c)\), then \(a \geq i\) which would either contradict the hook entry being the rightmost unpaired \(i\) or cell \((r, c+1)\) containing an \(i\). Thus, we only need to consider the subcases where the rightmost unpaired \(i\) is either in the leg or arm of \((r, c)\).
Assume that the rightmost unpaired $i$ in $A_T(r, c)$ for this entire paragraph. This implies that $\ell = i$. Since $(r, c + 1)$ contains an $i$, we have that $a < i$. If $\tilde{r} < r$, then $V_b$ acts on $T$ by removing $a$ from $(r, c)$, replacing $k$ with $a$ in $(\tilde{r}, c + 1)$, and appending $k$ to $A_T(\tilde{r}, c + 1)$. Since $a, k < i$, we have that $V_b$ does not change position of the rightmost unpaired $i$. Note that $(r + 1, c)$ still does not contain an $i + 1$ while $(r, c + 1)$ still contains an $i$. Thus, $f_i$ acts on $V_b(T)$ by removing the $i$ from $(r, c + 1)$ and appending an $i + 1$ to $L_{V_b(T)}(r, c)$. We now consider $V_b(f_i(T)).$ Note that $(r', c'), a', k'$ are the same as $(r, c), a,$ and $k$ respectively. Thus, $V_b$ acts in the same way as before. This gives us that $f_i(V_b(T)) = V_b(f_i(T)).$ If $\tilde{r} = r$, then $k$ is precisely the $i$ in cell $(r, c + 1)$. We see that $V_b$ acts on $T$ by removing $\{a, i\} \cap L_T(r, c)$ from $L_T(r, c)$ and $A_T(r, c)$, replacing $k$ with $\{(a, i) \cap L_T(r, c)\} \cup \{a\}$, and appending $k$ to $A_T(r + 1, c)$. Since there is an $i + 1$ in $L_{V_b(T)}(r, c + 1)$, we see that the rightmost unpaired $i$ in $V_b(T)$ is precisely $k$ in $A_{V_b(T)}(r, c + 1)$. Note that $(r + 1, c + 1)$ does not contain an $i + 1$ and $(r, c + 2)$ does not contain an $i$ because $(r, c + 1)$ contains an $i + 1$. Thus, $f_i$ acts on $V_b(T)$ by changing the $i$ in $A_{V_b(T)}(r, c + 1)$ to an $i + 1$. We now consider $V_b(f_i(T)).$ We have that $a'$ is the same as $a$ and $k'$ is the $i + 1$ in $(r, c + 1)$. We have $(a', \ell') \cap L_{f_i(T)}(r', c') = \{i + 1\} \cup \{(a, i) \cap L_T(r, c)\}$. This implies that $V_b$ acts on $f_i(T)$ by removing $\{i + 1\} \cup \{(a, i) \cap L_T(r, c)\} \cup \{a\}$ from $L_{f_i(T)}(r, c, c + 1)$, replacing $i + 1$ with $\{i + 1\} \cup \{(a, i) \cap L_T(r, c)\} \cup \{a\}$ in $(r, c, c + 1)$, and appending an $i + 1$ to $A_{f_i(T)}(r, c + 1).$ We see that $f_i(V_b(T)) = V_b(f_i(T)).$

Assume that the rightmost unpaired $i$ is in $A_T(r, c)$. This implies that $a = i$ and forces $a$ to correspond to the rightmost unpaired $i$. We also have that $k$ is the $i$ in $(r, c + 1)$. Since $i$ is weakly greater than all values in $(r, c)$, we have that $(a, \ell) \cap L_T(r, c) = \emptyset$. Thus, $V_b$ acts on $T$ by removing $a$ from $(r, c)$, replacing $k$ with $a$ in $(r, c + 1)$, and appending $k$ to $A_T(r, c + 1)$. Since $a$ was the rightmost unpaired $i$ in $T$ and cell $(r, c + 1)$ contains an $i + 1$ in its leg, we have that the rightmost unpaired $i$ in $V_b(T)$ is $k$ in $A_{V_b(T)}(r, c + 1)$. As $i + 1$ is in $(r, c + 1)$, we have that $(r + 1, c + 1)$ cannot contain an $i + 1$ and $(r, c + 2)$ cannot contain an $i$. This implies that $f_i$ acts on $V_b(T)$ by changing the $i$ in $A_{V_b(T)}(r, c + 1)$ to an $i + 1$. We now consider $V_b(f_i(T)).$ We have that $a'$ is the same as $a$ and $k'$ is equal to the $i + 1$ in $(r, c + 1)$. Note that $(a', \ell') \cap L_T(r, c) = \{i + 1\}$. This implies that $V_b$ acts on $f_i(T)$ by removing $i + 1$ from $L_{f_i(T)}(r, c)$ and $a$ from $A_{f_i(T)}(r, c)$, replacing the $i + 1$ with $\{i + 1, a\}$, and appending an $i + 1$ to $A_{f_i(T)}(r, c + 1)$. We see that $f_i(V_b(T)) = V_b(f_i(T)).$

• Assume that $(r + 1, c)$ does not contain an $i + 1$ and $(r, c + 1)$ does not contain an $i$. We break into subcases based on where the rightmost unpaired $i$ sits inside $(r, c)$.

Assume that the rightmost unpaired $i$ is in the hook entry of $(r, c)$ for the remainder of this paragraph. Note that this implies that $a > i$ and the rightmost unpaired $i$ in $V_b(T)$ is still the hook entry of $(r, c)$. We see that $V_b$ does not insert an $i + 1$ into $(r + 1, c)$ nor an $i$ into $(r, c + 1)$. This implies that $f_i$ acts on $T$ and $V_b(T)$ in the same way by changing the hook entry of $(r, c)$ into an $i + 1$. Next, we note that $(r', c'), a', k'$, and $(a', \ell) \cap L_{f_i(T)}(r', c')$ are the same as $(r, c), a, k$, and $(a, \ell) \cap L_T(r, c)$ respectively. Thus, $V_b$ acts on $T$ and $f_i(T)$ in the same manner without affecting the hook entry of $(r, c)$. Therefore, we have that the actions of $f_i$ and $V_b$ on $T$ are independent and $f_i(V_b(T)) = V_b(f_i(T)).$
Assume that the rightmost unpaired $i$ is in the leg of $(r, c)$ for the remainder of this paragraph. This implies that $a \neq i$. First, we assume that $a > i$ or $\tilde{r} < r$. Under this extra assumption, we observe that the action of $V_b$ does not change the position of the rightmost unpaired $i$. Also, $V_b$ does not insert an $i + 1$ into $(r + 1, c)$ nor an $i$ into $(r, c + 1)$. We see that $f_i$ acts on $T$ and $V_b(T)$ in the same way by changing the $i$ in the leg of $(r, c)$ into an $i + 1$. Next, we note that $(r', c')$, $a'$, and $k'$ are the same as $(r, c)$, $a$, and $k$ respectively. If $a > i$, we have that $a \geq i + 1$ implying that $(a', \ell') \cap L_{f_i(T)}((r', c')) = (a, \ell) \cap L_T(r, c)$. Thus, either $(a', \ell') \cap L_{f_i(T)}((r', c')) = (a, \ell) \cap L_T(r, c)$ or $\tilde{r} < r$. This implies that $V_b$ acts on $T$ and $f_i(T)$ in the same manner and does not affect the $i$ or $i + 1$ in the leg of $(r, c)$. Therefore, we have that the actions of $f_i$ and $V_b$ on $T$ are independent and $f_i(V_b(T)) = V_b(f_i(T))$. Next, assume that $\tilde{r} = r$ and $a < i$. This implies that $(a, \ell) \cap L_T(r, c) \neq \emptyset$ as $i \in (a, \ell) \cap L_T(r, c)$. We have that $V_b$ acts on $T$ by removing $(a, \ell) \cap L_T(r, c)$ from $L_T(r, c)$ and $a$ from $A_T(r, c)$, replacing $k$ with $(a, l) \cap L_T(r, c) \cup \{a\}$ in $(r, c + 1)$, and appending $k$ to $A_T(r, c + 1)$. By assumption, there was no $i$ in $(r, c + 1)$ to begin with. Thus, we have that the rightmost unpaired $i$ of $V_b(T)$ is the $i$ in $(r, c + 1)$ that replaced $k$. Since $k \geq i + 1$ or $k$ is the empty character, we have that the cell $(r + 1, c + 1)$ does not contain an $i + 1$ and the cell $(r, c + 2)$ does not contain an $i$. Hence, $f_i$ acts on $V_b(T)$ by replacing the $i$ in $L_{f_b(T)}((r, c + 1))$ with an $i + 1$. We now consider $V_b(f_i(T))$. We have that $f_i$ acts on $T$ by changing the $i$ in $L_T(r, c)$ to an $i + 1$. We see that $a'$ and $k'$ are the same as $a$ and $k$ respectively. Since $i > a$, we have that $i + 1 > a$ or in other words $i + 1 \in (a', \ell') \cap L_T(r, c)$. This implies that $(a', \ell') \cap L_{f_i(T)}((r', c')) = (((a', \ell') \cap L_T(r, c)) \cup \{i + 1\}) - \{i\}$. We have $V_b$ acts on $f_i(T)$ by removing $(a', \ell') \cap L_{f_i(T)}((r, c))$ from $L_{f_i(T)}((r, c))$ and $a$ from $A_{f_i(T)}((r, c))$, replacing $k$ with $(a', \ell') \cap L_{f_i(T)}((r, c))$ in $(r, c + 1)$, and appending $k$ to $A_{f_i(T)}((r, c + 1))$. We see that $f_i(V_b(T)) = V_b(f_i(T))$. Assume that the rightmost unpaired $i$ is in $A_T(r, c)$ and $\tilde{r} < r$ or $(a, \ell) \cap L_T(r, c) = \emptyset$ for this entire paragraph. Under this assumption, $f_i$ acts on $T$ by changing the rightmost $i$ in the arm of $(r, c)$ to an $i + 1$. Also, $V_b$ acts on $T$ by removing $a$ from $A_T(r, c)$, replacing $k$ in $(\tilde{r}, c + 1)$ with $a$, and appending $k$ to $A_T(\tilde{r}, c + 1)$. First, we make the additional assumption that $i < a$. Since we assume the rightmost unpaired $i$ is in the arm of $(r, c)$ and $i < a$, we have the rightmost unpaired $i$ in $V_b(T)$ is in the same position as in $T$. Note that the cell $(r + 1, c)$ still does not contain an $i + 1$ and the cell $(r, c + 1)$ still does not contain an $i$. Thus, we have that $f_i$ acts on $V_b(T)$ by changing the rightmost $i$ in $A_{V_b}(r, c)$ into an $i + 1$. We now consider $V_b(f_i(T))$. We see that $a'$ and $k'$ are the same as $a$ and $k$ respectively. This implies that $V_b$ acts on $f_i(T)$ by removing $a$ from $(r, c)$, replacing $k$ with $a$ in $(\tilde{r}, c)$, and appending $k$ to $A_{f_i(T)}(\tilde{r}, c + 1)$. We see that $f_i(V_b(T)) = V_b(f_i(T))$. Next, we make the assumption that $a = i$ and column $c + 1$ does not contain both an $i$ and an $i + 1$. We have that the rightmost unpaired $i$ in $V_b(T)$ is precisely the $i$ that replaced $k$ in $(\tilde{r}, c + 1)$. We also have that $k \geq i + 1$ or $k$ is the empty character implying that the cell $(\tilde{r} + 1, c + 1)$ does not contain an $i + 1$ and the cell $(\tilde{r}, c + 2)$ does not contain an $i$. This implies that $f_i$ acts on $V_b(T)$ by changing the $i$ in $L_{V_b(T)}^+(\tilde{r}, c + 1)$ to an $i + 1$. We now consider $V_b(f_i(T))$. We see that $a'$ is the $i + 1$ in $(r, c)$ created by appending $f_i$ and $k'$ is the same as $k$. Thus, $V_b$ acts on $f_i(T)$ by removing the $i + 1$ from $(r, c)$, replacing $k$ with an $i + 1$ in $(\tilde{r}, c)$, and appending $k$ to $A_{f_i(T)}(\tilde{r}, c + 1)$. We see that $f_i(V_b(T)) = V_b(f_i(T))$. Next, we assume that $a = i$ and column $c + 1$ contains both an $i$ and an $i + 1$ in the same cell. Note that this implies that $k = i$. Since $a$ corresponded to the rightmost unpaired $i$ in $T$ and the only cell in column $c + 1$ that contains an $i + 1$ or an $i$ is $(\tilde{r}, c + 1)$, we have
that the rightmost unpaired $i$ in $V_b(T)$ corresponds to the $i$ appended to $\mathcal{A}_T(\tilde{r}, c + 1)$. Since $(\tilde{r}, c + 1)$ contains an $i + 1$ in $V_b(T)$, we have that $(\tilde{r} + 1, c + 1)$ cannot contain an $i + 1$ and $(\tilde{r}, c + 2)$ cannot contain an $i$. Thus, $f_i$ acts on $V_b(T)$ by changing the $i$ in $A_{V_b(T)}(\tilde{r}, c + 1)$ to an $i + 1$. We now consider $V_b(f_i(T))$. We see that $a'$ is the $i + 1$ in $(r, c)$ obtained after applying $f_i$ and $k'$ is the $i + 1$ in cell $(\tilde{r}, c + 1)$. This implies that $V_b$ acts on $f_i(T)$ by removing the $i + 1$ from $(r, c)$, replacing $k'$ with an $i + 1$ in $(\tilde{r}, c + 1)$, and appending $k'$ to $A_{f_i(T)}(\tilde{r}, c + 1)$. We see that $f_i(V_b(T)) = V_b(f_i(T))$. Finally, we make the assumption that $a = i$ and column $c + 1$ contains both an $i$ and an $i + 1$ but in different cells. We once again have that $k = i$, but now we have that $(\tilde{r} + 1, c + 1)$ contains an $i + 1$. We have that the rightmost unpaired $i$ in $V_b(T)$ is the $i$ that was appended to $\mathcal{A}_T(\tilde{r}, c + 1)$. Since $(\tilde{r} + 1, c + 1)$ contains an $i + 1$, we have that $f_i$ acts on $V_b(T)$ by removing the $i$ from $A_{V_b(T)}(\tilde{r}, c + 1)$ and appending an $i + 1$ to $A_{V_b(T)}(\tilde{r} + 1, c + 1)$. We now consider $V_b(f_i(T))$. We see that $a'$ is the $i + 1$ in $(r, c)$ obtained after applying $f_i$ and $k'$ the $i + 1$ in cell $(\tilde{r} + 1, c + 1)$. This implies that $V_b$ acts on $f_i(T)$ by removing the $i + 1$ from $(r, c)$, replacing $k'$ with an $i + 1$ in $(\tilde{r} + 1, c + 1)$, and appending $k'$ to $A_{f_i(T)}(\tilde{r} + 1, c + 1)$. We see that $f_i(V_b(T)) = V_b(f_i(T))$.

Assume that the rightmost unpaired $i$ is in the arm of $(r, c)$, $\tilde{r} = r$, and $(a, \ell] \cap L_T(r, c) \neq \emptyset$ for this entire paragraph. First, we make the additional assumption that $i < a$. This gives us that $V_b(T)$ is attained from $T$ by removing $(a, \ell] \cap L_T(r, c)$ from $L_T(r, c)$ and $a$ from $\mathcal{A}_T(r, c)$, replacing $k$ in cell $(r, c + 1)$ with $((a, \ell] \cap L_T(r, c)) \cup \{a\}$, and appending $k$ to $\mathcal{A}_T(r, c + 1)$. Since $k, a > i$, we have that the rightmost unpaired $i$ in $V_b(T)$ remains the same as in $T$. We also have that the cell $(r + 1, c)$ does not contain an $i + 1$ and the cell $(r, c + 1)$ does not contain an $i$. Thus, $f_i$ acts on $V_b(T)$ by changing the rightmost $i$ in $A_{V_b(T)}(r, c)$ to an $i + 1$. We now consider $V_b(f_i(T))$. We have that $f_i$ acts on $T$ by changing the rightmost $i$ in $\mathcal{A}_T(r, c)$ to an $i + 1$. We see that $a', k'$, and $(a', \ell'] \cap L_{f_i(T)}(r', c')$ are the same as $a, k$, and $(a, \ell] \cap L_T(r, c)$ respectively. This implies that $V_b$ acts on $f_i(T)$ by removing $(a, \ell] \cap L_T(r, c)$ from $L_{f_i(T)}(r, c)$ and $a$ from $\mathcal{A}_{f_i(T)}(r, c)$, replacing $k$ in cell $(r, c + 1)$ with $((a, \ell] \cap L_T(r, c)) \cup \{a\}$, and appending $k$ to $\mathcal{A}_{f_i(T)}(r, c + 1)$. We see that $f_i(V_b(T)) = V_b(f_i(T))$. Next, we assume that $a = i$ and $(r, c)$ contains an $i + 1$. Since $a = i$, the $i + 1$ in $(r, c)$ must be in its leg. Also as $a$ is the rightmost unpaired $i$ of $T$, we must have that $(r, c)$ contains another $i$ besides $a$. This gives us that $V_b(T)$ is attained from $T$ by removing $(a, \ell] \cap L_T(r, c)$ from $L_T(r, c)$ and $a$ from $\mathcal{A}_T(r, c)$, replacing $k$ in cell $(r, c + 1)$ with $((a, \ell] \cap L_T(r, c)) \cup \{a\}$, and appending $k$ to $\mathcal{A}_T(r, c + 1)$. Note that the $i$ inserted into $(r, c + 1)$ becomes $i$-paired while an $i$ in $(r, c)$ becomes unpaired. This implies that the rightmost unpaired $i$ in $V_b(T)$ still sits in the cell $(r, c)$. We see that the cell $(r + 1, c)$ still does not contain an $i + 1$; however, the cell $(r, c + 1)$ now contains an $i$. This implies that $f_i$ acts on $V_b(T)$ by removing the $i$ from the cell $(r, c + 1)$ and appending an $i + 1$ to $L_{V_b(T)}(r, c)$. We now consider $V_b(f_i(T))$. We have that $f_i$ acts on $T$ by changing $a$ into an $i + 1$. We have that $a'$ is the $i + 1$ obtained from applying $f_i$ and $k'$ is the same as $k$. We see that $(a', \ell'] \cap L_{f_i(T)}(r', c')$ is the same as $(a, \ell] \cap L_T(r, c)$ excluding the $i + 1$. We have that $V_b$ acts on $f_i(T)$ by removing $(a', \ell'] \cap L_{f_i(T)}(r', c')$ from $L_{f_i(T)}(r, c)$ and $i + 1$ from $\mathcal{A}_{f_i(T)}(r, c)$, leaving the $i + 1$ in $L_{f_i(T)}(r, c)$, replacing $k$ in $(r, c + 1)$ with $((a', \ell] \cap L_{f_i(T)}(r', c')) \cup \{a\}$, and appending $k$ to $\mathcal{A}_{f_i(T)}(r, c + 1)$. We see that $f_i(V_b(T)) = V_b(f_i(T))$. Finally, we assume that $a = i$ and $i + 1$ is not in the cell $(r, c)$. This gives us that $V_b(T)$ is attained from $T$ by removing $(a, \ell] \cap L_T(r, c)$ from $L_T(r, c)$ and $a$ from $\mathcal{A}_T(r, c)$, replacing $k$ in cell $(r, c + 1)$ with $((a, \ell] \cap L_T(r, c)) \cup \{a\}$, and appending $k$ to $\mathcal{A}_T(r, c + 1)$. Since $k \geq j > i + 1$ for all $j \in (a, \ell] \cap L_T(r, c)$, we have that the $i$ inserted into the cell
(r, c + 1) is the rightmost unpaired i in V_b(T). Note that the cell (r + 1, c + 1) does not contain an i + 1 and the cell (r, c + 2) does not contain an i. Thus, f_i acts on V_b(T) by changing the i in (r, c + 1) to an i + 1. We now consider V_b(f_i(T)). We have that f_i acts on T by changing a into an i + 1. We have that a' is the i + 1 obtained from applying f_i and k' is the same as k. We see that \((a', \ell'] \cap L_{f_i(T)}(r', c') = (a, \ell] \cap L_T(r, c)).\]

We have that V_b acts on f_i(T) by removing \((a, \ell] \cap L_T(r, c)) from L_{f_i(T)}(r, c) and i + 1 from A_{f_i(T)}(r, c), replacing k in (r, c + 1) with \(((a, \ell] \cap L_T(r, c)) \cup \{a', \ell'] \}, and appending k to A_{f_i(T)}(r, c + 1). We see that f_i(V_b(T)) = V_b(f_i(T)).

Case 2: Assume that \(\hat{r} < r\) and \(\hat{c} = c\).

Note that \(a > i\). By Lemma 3.10 we have that \(R_i(T) = R_i(V_b(T))\) unless \(a = i + 1\) and column \(c + 1\) contains both an i and an i + 1. However, even in this special case, we see that the rightmost unpaired i of V_b(T) is in the same position as the rightmost unpaired i of T. We also see that V_b(T) does not change whether or not cell \((\hat{r} + 1, c)\) contains an i + 1 and whether or not cell \((\hat{r}, c + 1)\) contains an i. Thus, f_i acts on the same i and in the same way for both T and V_b(T). Since \(a > i\), we have that k' is the same as k. Note that the only way for f_i to affect the cell \((r, c)\) in T is if \(\hat{r} = r - 1\) and \((r, c)\) contains an i + 1. However, even in this special case, we see that \((r', c'), a', \ell', (a', \ell'] \cap L_{f_i(T)}(r', c')\) are the same as \((r, c), a, \ell, (a, \ell] \cap L_T(r, c))\). Thus, V_b acts on T and f_i(T) in the same way. Therefore, we have that the actions of f_i and V_b on T are independent and f_i(V_b(T)) = V_b(f_i(T)).

Case 3: Assume that \(\hat{c} < c\).

Let \(\hat{i}\) denote the rightmost unpaired i of T. From the proof of Lemma 3.10, we have that V_b does not change whether or not the i's to the right of \(\hat{i}\) in R_i(T) are i-paired. Thus, the rightmost unpaired i in R_i(T) and R_i(V_b(T)) are in the same position. As V_b does not affect any column to the left of column c, we have that the rightmost unpaired i for V_b(T) is in the same position as the rightmost unpaired i for T. Note that V_b also does not affect whether or not cell \((\hat{r} + 1, \hat{c})\) contains an i + 1 and whether or not cell \((\hat{r}, \hat{c} + 1)\) contains an i. Thus, f_i acts on the rightmost unpaired i in T and V_b(T) in exactly the same way. Next, we note that \((\hat{r}', c'), a', \ell', (a', \ell'] \cap L_{f_i(T)}(r', c')\) are the same as \((r, c), a, \ell, k, (a, \ell] \cap L_T(r, c))\) respectively. Thus, V_b acts on T and f_i(T) in the same way. Therefore, we have that the actions of f_i and V_b on T are independent and f_i(V_b(T)) = V_b(f_i(T)).

Case 4: Assume that \(\hat{r} \leq r\) and \(\hat{c} = c + 1\).

Under this assumption, we have that column \(c + 1\) does not contain an i + 1 and \(a \neq i + 1\) since the cells in column \(c + 1\) do not contain any arms. We break into subcases.

- Assume that \(k \neq i\). This implies that the rightmost unpaired i in V_b(T) is in the same position as the rightmost unpaired i in T. We see that V_b does not change whether or not cell \((\hat{r} + 1, c + 1)\) contains an i + 1 and whether or not cell \((\hat{r}, c + 2)\) contains an i. Thus, f_i acts on the rightmost unpaired i in T and V_b(T) in exactly the same way. We also observe that \((\hat{r}', c'), a', \ell', k', (a', \ell'] \cap L_{f_i(T)}(r', c')\) are the same as \(a, \ell, k, (a, \ell] \cap L_T(r, c))\) respectively. Thus, V_b acts on T and f_i(T) in the same way. Therefore, we have that the actions of f_i and V_b on T are independent and f_i(V_b(T)) = V_b(f_i(T)).

- Assume that \(k = i\). We see that the rightmost unpaired i in V_b(T) is the i that was appended to A_T(\(\hat{r}, c + 1\)). Note that V_b does not change whether or not cell \((\hat{r} + 1, c + 1)\) contains an i + 1 and whether or not cell \((\hat{r}, c + 2)\) contains an i. We first make the extra assumption that \((\hat{r}, c + 2)\) in T contains an i. This implies that f_i acts on V_b(T) and T in the same way by removing the i from the hook entry of \((\hat{r}, c + 2)\) and appending an i + 1 to the leg of \((\hat{r}, c + 1)\). We also have that \((\hat{r}', c'), a', \ell', k', (a', \ell'] \cap L_{f_i(T)}(r', c')\) are equal to \((r, c), a, \ell, k, (a, \ell] \cap L_T(r, c))\) respectively. Thus, V_b acts on T and f_i(T) in the same way. Therefore, we have that the actions of f_i and V_b on T are independent and f_i(V_b(T)) = V_b(f_i(T)). We now assume that \((\hat{r}, c + 2)\) does not contain an i. This
Proof. In Theorem 3.14, let \( U \) be the uncrowding map and denote the pair of insertion and recording tableaux produced at the \( i \)-th step for \( 0 \leq i \leq \alpha \) of the uncrowding map \( U \) for \( T \) and \( f_i(T) \) as \( (P_j(T), Q_j(T)) \) and \( (P_j(f_i(T)), Q_j(f_i(T))) \), respectively. As crystal operators do not change the shape of \( T \), we have \( \text{shape}(P_j(T)) = \text{shape}(f_i(T)) \) for all \( 0 \leq j \leq \alpha \). Hence

\[
\text{shape}(P_{j+1}(T))/\text{shape}(P_j(T)) = \text{shape}(P_{j+1}(f_i(T)))/\text{shape}(P_j(f_i(T))) \quad \text{for all } 0 \leq j \leq \alpha - 1.
\]

Next we show \( Q_j(T) = Q_j(f_i(T)) \) for all \( 0 \leq j \leq \alpha \) by induction. When \( j = 0 \), \( Q_0(T) = Q_0(f_i(T)) \) since \( \text{shape}(R_0(T)) = \text{shape}(R_0(f_i(T))) = \text{shape}(T) \).

Suppose \( Q_j(T) = Q_j(f_i(T)) \) for a given \( j \geq 0 \). It suffices to show that the cells

\[
\text{shape}(Q_{j+1}(T))/\text{shape}(Q_j(T)) = \text{shape}(P_{j+1}(T))/\text{shape}(P_j(T)) \quad \text{and} \quad \text{shape}(Q_{j+1}(f_i(T)))/\text{shape}(Q_j(f_i(T))) = \text{shape}(P_{j+1}(f_i(T)))/\text{shape}(P_j(f_i(T)))
\]

in \( Q_{j+1}(T) \) and \( Q_{j+1}(f_i(T)) \) are at the same position with the same entry. By (3.2), the cells are in the same position, say in column \( \hat{c} \). By Definition 2.5, \( f_i \) does not move elements in the arm.
to a different column, so the columns in which we start the uncrowding insertion \( V \) on \( P_j(T) \) and \( P_j(f_i(T)) \) are the same, say \( c \), by Definition 3.5. Hence the cells \( \text{shape}(Q_{j+1}(T))/\text{shape}(Q_j(T)) \) and \( \text{shape}(Q_{j+1}(f_i(T)))/\text{shape}(Q_j(f_i(T))) \) are at the same position with entry \( \tilde{c} - c \). The theorem follows. \( \square \)

Hawkes and Scrimshaw [HS20, Theorem 4.6] proved that \( \text{HVT}^m(\lambda) \) is a Stembridge crystal by checking the Stembridge axioms. This also follows directly from our analysis above.

**Corollary 3.15.** The crystal \( \text{HVT}^m(\lambda) \) of Definition 2.5 is a Stembridge crystal of type \( A_{m-1} \).

**Proof.** According to [MPS21], \( \text{SVT}^m(\mu) \) is a Stembridge crystal of type \( A_{m-1} \). By Theorem 3.14, the map

\[
\mathcal{U} : \text{HVT}^m(\lambda) \rightarrow \bigsqcup_{\mu \supseteq \lambda} \text{SVT}^m(\mu) \times \hat{\mathcal{F}}(\mu/\lambda),
\]

is a strict crystal morphism (see for example [BS17, Chapter 2]). The statement follows. \( \square \)

### 3.4. Uncrowding map on multiset-valued tableaux

The uncrowding map on hook-valued tableaux described above turns out to be a generalization of the uncrowding map on multiset-valued tableaux by Hawkes and Scrimshaw [HS20, Section 3.2]. We will prove that this is indeed the case in this section. Let us recall the definition of the uncrowding map in [HS20, Section 3.2].

**Definition 3.16.** Let \( T \in \text{MVT}(\lambda) \). The uncrowding map

\[
\Upsilon : \text{MVT}(\lambda) \rightarrow \bigsqcup_{\mu \supseteq \lambda} \text{SSYT}(\mu) \times \hat{\mathcal{F}}(\mu/\lambda)
\]

sends \( T \) to a pair of tableaux using the following algorithm:

1. Set \( U_{\lambda_1+1} = \emptyset \) and \( F_{\lambda_1+1} \) be the unique column-flagged increasing tableau of shape \( \emptyset/\emptyset \).

2. Let \( 1 \leq k \leq \lambda_1 \) and assume that the pair \( (U_{k+1}, F_{k+1}) \) is defined. The pair \( (U_k, F_k) \) is defined recursively from \( (U_{k+1}, F_{k+1}) \) using the following two steps:
   (a) Define \( U_k \) as the RSK row insertion tableau from the word

   \[
   R(C_k)R(C_{k+1}) \cdots R(C_{\lambda_1}),
   \]

   where \( C_j \) is the \( j \)-th column of \( T \) for every \( 1 \leq j \leq \lambda_1 \). In other words, if we denote by \( T_{\geq k} \) the tableau formed by the columns weakly to the right of the \( k \)-th column of \( T \), \( U_k \) is obtained by performing the RSK row insertion using the column reading word of \( T_{\geq k} \).
   (b) Form the tableau \( F_k \) of shape \( \text{shape}(U_k)/\text{shape}(T_{\geq k}) \) as follows. Shift \( F_{k+1} \) by one column to the right and fill the boxes in the same positions into \( F_k \); for every unfilled box in the shape \( \text{shape}(U_k)/\text{shape}(U_{k+1}) \), label each box in column \( i \) with entry \( i - 1 \).

Define \( \Upsilon(T) = (U, F) := (U_1, F_1) \).

**Example 3.17.** Let \( T \) be the multiset-valued tableau

\[
T = \begin{bmatrix}
45 \\
233 & 345 \\
1 & 11 & 4
\end{bmatrix}
\]
Then, we obtain the following pairs of tableaux for the uncrowding map $\Upsilon$:

$$
(U_1, F_1) = (\emptyset, \emptyset) \\
(U_2, F_2) = \begin{pmatrix}
3 & 5 \\
1 & 1 & 4 \\
1 & 1 & 3 & 4 & 4 \\
\end{pmatrix}, \ 
\begin{pmatrix}
1 & 1 \\
2 & 3 \\
2 & 3 & 5 \\
\end{pmatrix} = (U, F) = \Upsilon(T).
$$

**Proposition 3.18.** Let $T \in \text{MVT}(\lambda)$. Then $U(T) = \Upsilon(T)$. In other words, the uncrowding map as defined in Definition 3.5 is equivalent to the uncrowding map of Definition 3.16 in [HS20, Section 3.2].

**Proof.** Recall from Definition 3.5, that the pair of uncrowding and recording tableaux for $U(T)$ is denoted by $(P(T), Q(T)) = U(T)$. Similarly, let us denote $(U(T), F(T)) := \Upsilon(T)$.

Assume that $S \in \text{MVT}(\lambda)$ is highest weight, that is, $e_i(S) = 0$ for $i \geq 1$. By [HS20, Proposition 3.10], row $i$ of $S$ only contains the letter $i$. Thus its weight is some partition $\mu = (\mu_1, \mu_2, \ldots, \mu_{\ell})$ if $\lambda = (\lambda_1, \lambda_2, \ldots, \lambda_{\ell})$. By Proposition 3.12 and Theorem 3.14, $P(S) \in \text{SSYT}$ is highest weight. As weights of tableaux are preserved under uncrowding, the weight of $P(S)$ is equal to $\mu$. By a similar argument using [HS20, Theorem 3.17], $U(S) \in \text{SSYT}$ is also highest weight with weight $\mu$. Since highest weight semistandard Young tableaux are uniquely determined by their weights, we have $P(S) = U(S)$.

Recall that as long as $f_iT \neq 0$ for $T \in \text{MVT}(\lambda)$, we have $U(f_iT) = f_iU(T)$ by [HS20, Theorem 3.17] and $P(f_iT) = f_iP(T)$ by Theorem 3.14. Now let $T \in \text{MVT}(\lambda)$ be arbitrary. Then $T = f_{i_1} \cdots f_{i_k}S$ for some sequence of $i_1, \ldots, i_k$ and $S$ highest weight. Hence,

$$
P(T) = P(f_{i_1} \cdots f_{i_k}S) = f_{i_1} \cdots f_{i_k}P(S) = f_{i_1} \cdots f_{i_k}U(S) = U(f_{i_1} \cdots f_{i_k}S) = U(T).
$$

It remains to show that $Q(T) = F(T)$ for all $T \in \text{MVT}(\lambda)$. To do this, we show that the newly created boxes of the uncrowding map up to a specified column in Definition 3.16 are in the same positions as those for the uncrowding insertion in Definition 3.5. For every $Y \in \text{MVT}(\mu)$ and for every $1 \leq j \leq \mu_1$, denote by $Y_{\geq j}$ the tableau formed by the rightmost $j$ columns of $Y$; here $Y_{\geq \mu_1+1}$ is the empty tableau.

Let $T \in \text{MVT}(\lambda)$ be arbitrary. For $1 \leq k \leq \lambda_1 + 1$, let $P^{(k)}$ be the tableau obtained by performing the uncrowding map $U$ on $T$ on the columns from right to left up to and including the $k$-th column of $T$; here $P^{(\lambda_1+1)} = T$. In other words, $P^{(k)} = Y^{\alpha_k}(T)$ as in Definition 3.4, where $\alpha_k$ is the arm excess of $T_{\geq k}$. As the entries to the left of column $k$ of $T$ are untouched by the uncrowding insertion in Definition 3.4, for every $1 \leq k \leq \lambda_1 + 1$, we have $(P^{(k)})_{\geq k} = P(T_{\geq k}) = U(T_{\geq k})$. It follows that for every $1 \leq k \leq \lambda_1$, up to horizontal shifts, the newly formed boxes in $\text{shape}(P^{(k)})/\text{shape}(P^{(k+1)}) = \text{shape}[(P^{(k)})_{\geq k+1}]/\text{shape}[(P^{(k+1)})_{\geq k+1}]$ and $\text{shape}[(U(T_{\geq k}))_{\geq k+1}]/\text{shape}[(U(T_{\geq k+1}))_{\geq k+1}]$ are in the same positions. Since the entries in these boxes both record the difference in column indices relative to the $k$-th column for each $1 \leq k \leq \lambda_1$ and since the recording tableaux for both maps are formed from the union of these boxes, we conclude that $Q(T) = F(T)$, completing the proof.  

3.5. Crowding map. In this section, we give a description of the “inverse” of the uncrowding map.
We begin by introducing some notation. Let \( F \in \hat{\mathcal{F}} \) with \( e \) entries. For each cell \((r, c)\) in \( F\) with entry \( F(r, c)\), define the corresponding \textit{destination column} to be \( d(r, c) = c - F(r, c)\). Define the \textit{crowding order} on \( F\) by ordering all the cells in \( F\) with a filling, first determined by their destination column (smallest to largest) and then by column index (largest to smallest). Denote the order by \((r_1, c_1), (r_2, c_2), \ldots, (r_e, c_e)\). Set \( \alpha(F) = (\alpha_1, \alpha_2, \ldots, \alpha_e)\), where \( \alpha_i = F(r_i, c_i)\). Let the arm excess for a column of a hook-valued tableau be the sum of arm excesses of all its cells.

**Definition 3.19.** Let \( h \in \text{HVT} \) and let \((r, c)\) be a cell in \( h\) with \( c > 1\) and with at most one element in \( A_h(r, c)\). If \( A_h(r, c)\) is empty, we also require that the cell \((r, c)\) is a corner cell in \( h\). Then we define the crowding \( C_h\) on the pair \([h, (r, c)]\) by the following algorithm:

1. If \( A_h(r, c)\) is nonempty, set \( m \) to be the only element in \( A_h(r, c)\) and \( b = \max\{x \in L_h^+(r, c) \mid x \leq m\} \). Otherwise, set \( m = H_h(r, c)\) and \( b = \max\{L_h^+(r, c)\} \).
2. Find the largest \( r'\) such that \( H_h(r', c - 1) \leq b\). If \( r' = r\), set \( q = H_h(r, c)\). Otherwise, set \( q = b\).
3. In either case, append \( q\) to \( A_h(r', c - 1)\).

(a) If \( r'\) from Step 2 equals \( r\), perform either of the following:

(i) If \( A_h(r, c)\) is nonempty, move the set \( \{x \in L_h(r, c) \mid q < x \leq m\} \) from \( L_h(r, c)\) to \( L_h(r', c - 1)\) and keep it strictly increasing. Remove \( m\) from \( A_h(r, c)\) and set \( H_h(r, c) = m\).

(ii) Otherwise, \( A_h(r, c)\) is empty, so move \( L_h(r, c)\) into \( L_h(r', c - 1)\) and keep it to be strictly increasing. Remove cell \((r, c)\) from \( h\).

(b) Otherwise, \( r' \neq r\) and perform either of the following:

(i) Suppose that \( A_h(r, c)\) is nonempty. Replace \( q\) in \( L_h^+(r, c)\) with \( m\). Remove \( m\) from \( A_h(r, c)\).

(ii) If instead \( A_h(r, c)\) is empty, then remove cell \((r, c)\) from \( h\).

Denote the resulting (not necessarily semistandard) hook-valued tableau by \( h'\). We write \( C_h([h, (r, c)]) = [h', (r', c - 1)]\). We also define the projections \( p_1\) and \( p_2\) by \( p_1 \circ C_h([h, (r, c)]) = h'\) and \( p_2 \circ C_h([h, (r, c)]) = (r', c - 1)\). See Figures 3 and 4 for illustration.

**Figure 3.** When \( r' = r\). Left: (i) \( A_h(r, c) \neq \emptyset\). Right: (ii) \( A_h(r, c) = \emptyset\).

**Figure 4.** When \( r' \neq r\). Left: \( A_h(r, c) \neq \emptyset\). Right: \( A_h(r, c) = \emptyset\).

**Example 3.20.** We compute \( C_h\) in two examples:

\[
T = \begin{bmatrix} 5 & 5 & 4 & 3 \\ 4 & 3 & 1 & 1 \\ 3 & 2 & 4 & 1 \\ 1 & 2 & 4 & 1 \end{bmatrix}, \quad C_h([T, (1, 2)]) = \begin{bmatrix} 5 & 5 & 4 & 3 \\ 4 & 3 & 1 & 1 \\ 3 & 2 & 4 & 1 \\ 1 & 2 & 4 & 1 \end{bmatrix} = [T', (1, 1)].
\]
With the same notation as above, define the tableaux

\[ T \]

Remark 3.21. In Definition 3.19,

- if \( r' = r \), then \( h' \) is always semistandard and has the same weight as \( h \);
- if \( r' \neq r \) and \( A_h(r, c) \) is empty, then \( h' \) might have fewer letters than \( h \). In Example 3.20, \( S \) contains 5 letters while \( S' \) only contains 4. This happens precisely when \( L_h(r, c) \) is nonempty.

In principle, the arm in cell \((r', c - 1)\) could be greater than the \( q \) that is to be inserted. However, we only consider the cases as defined in the order described by the next paragraph. We refer to Proposition 3.27 which states that all tableaux we deal with in this section are indeed semistandard hook-valued tableaux.

Let \((S, F) \in \text{SVT}(\mu) \times \hat{F}(\mu/\lambda)\) with crowding order \((r_1, c_1), (r_2, c_2), \ldots, (r_e, c_e)\) and \(\alpha(F) = (\alpha_1, \alpha_2, \ldots, \alpha_e)\). For all \(0 \leq j \leq e - 1\) and for all \(0 \leq s \leq \alpha_{j+1}\), define \(T_j^{(s)}\) recursively by setting \(T_0^{(0)} := S\) and

\[
T_j^{(s)} := \begin{cases} 
p_1 \circ C_b([T_{j-1}^{(s-1)}, (r_j+1, c_j+1)]) & \text{when } s > 0, \\
T_{j-1}^{(s)} & \text{when } s = 0 \text{ and } j > 0.
\end{cases}
\]

Additionally, define \(T_e^{(0)} := T_e^{(\alpha_e)}\).

Thus we obtain the following sequence

\[
S = T_0^{(0)} \xrightarrow{p_1 \circ C_b^{(r_1,c_1)}} T_1^{(0)} \xrightarrow{p_1 \circ C_b^{(r_2,c_2)}} T_2^{(0)} \xrightarrow{p_1 \circ C_b^{(r_3,c_3)}} \ldots \xrightarrow{p_1 \circ C_b^{(r_e,c_e)}} T_e^{(0)}.
\]

Remark 3.22. The tableaux \(T_j^{(s)}\) are well-defined. We check the conditions in Definition 3.19. Let \(h = T_j^{(s)}\) for some \(0 \leq j \leq e - 1\) and for some \(0 \leq s < \alpha_{j+1}\), with cell \((r, c)\).

- Since \(F \in \hat{F}\), we always have \(c > 1\).
- The case that \(A_h(r, c)\) is empty can only occur in \(T_{j-1}^{(0)}\) for some \(j > 0\). In this case, \((r, c) = (r_j, c_j)\), which is a corner cell.
- Consider the \(\alpha_j\) steps in \(T_{j-1}^{(0)} \xrightarrow{p_1 \circ C_b^{(r_j,c_j)}} T_j^{(0)}\). We first delete cell \((r_j, c_j)\), which has no arm.

Then at every step after that, we move leftward one column at a time. Before we reach column \(d(r_j, c_j)\), there is exactly one column with arm excess being 1 and the rest has zero arm excess among columns to the right of \(d(r_j, c_j)\) since recall that the cells \((r_j, c_j)\) are ordered from smallest to largest destination column. Once we reach column \(d(r_j, c_j)\), the cell there may contain more than one arm element, but we then go to \((r_{j+1}, c_{j+1})\), which is a corner cell instead. Thus there is at most one element in \(A_h(r, c)\).

Definition 3.23. With the same notation as above, define the insertion path of \(T_{j-1}^{(0)} \rightarrow T_j^{(0)}\) for \(1 \leq j \leq e\) to be

\[
\text{path}_j := \left((r_j^{(0)}, c_j^{(0)}), (r_j^{(1)}, c_j^{(1)}), \ldots, (r_j^{(\alpha_j)}, c_j^{(\alpha_j)})\right),
\]

where \((r_j^{(s)}, c_j^{(s)}) := p_2 \circ C_b([T_{j-1}^{(0)}, (r_j, c_j)])\) for \(0 \leq s \leq \alpha_j\).
Example 3.24. Consider the following pair of tableaux \((S, F) \in \text{HVT}((5, 3, 2)) \times \hat{\mathcal{F}}((5, 3, 2)/((3, 2, 1)))\),

\[
S = \begin{array}{ccc}
5 & 4 & 5 \\
4 & 3 & 3 \\
1 & 2 & 4 \\
\end{array}, \quad F = \begin{array}{ccc}
1 &  &  \\
 &  &  \\
1 & 3 & 4 \\
\end{array}
\]

The crowding order is \((1, 5), (1, 4), (3, 2), (2, 3)\). The insertion path and destination column for each of them are:

\[
\text{path}_1 = ((1, 5), (1, 4), (2, 3), (2, 2), (2, 1)), \quad d(1, 5) = 1,
\]

\[
\text{path}_2 = ((1, 4), (2, 3), (2, 2), (3, 1)), \quad d(1, 4) = 1,
\]

\[
\text{path}_3 = ((3, 2), (3, 1)), \quad d(3, 2) = 1,
\]

\[
\text{path}_4 = ((2, 3), (2, 2)), \quad d(2, 3) = 2.
\]

We obtain the sequence from the algorithm:

\[
\begin{array}{cccccc}
5 & 4 & 5 & 5 & 4 & 5 \\
4 & 3 & 3 & 23 & 3 & 4 \\
1 & 2 & 4 & 1 & 1 & 4 \\
\end{array}, \quad p_1 \circ C_{(1,5)}^{\text{a}} = (1, 5) \\
\begin{array}{cccccc}
5 & 4 & 5 & 5 & 4 & 5 \\
23 & 3 & 4 & 23 & 3 & 4 \\
1 & 1 & 4 & 1 & 1 & 4 \\
\end{array}, \quad p_1 \circ C_{(1,4)}^{\text{a}} = (1, 4) \\
\begin{array}{cccccc}
5 & 44 & 5 & 5 & 44 & 5 \\
5 & 44 & 5 & 23 & 3 & 4 \\
1 & 1 & 2 & 1 & 1 & 2 \\
\end{array}, \quad p_1 \circ C_{(3,2)}^{\text{a}} = (3, 2) \\
\begin{array}{cccccc}
5 & 44 & 5 & 5 & 44 & 5 \\
23 & 4 & 3 & 4 & 23 & 4 \\
1 & 1 & 1 & 1 & 1 & 1 \\
\end{array}, \quad p_1 \circ C_{(2,3)}^{\text{a}} = (2, 3)
\]

Lemma 3.25. If \(d(r_j, c_j) = d(r_{j+1}, c_{j+1})\), then \(\text{path}_{j+1}\) is weakly above \(\text{path}_j\).

Proof. By the definition of crowding order, \(d(r_j, c_j) = d(r_{j+1}, c_{j+1})\) implies \(c_j > c_{j+1}\). Set \(z_j := c_j - c_{j+1}\). Then we have \(c_j^{(s+z_j)} = c_j - z_j - s = c_{j+1} - s = c_{j+1}^{(s)}\) for \(0 \leq s \leq \alpha_{j+1}\). We need to show that \(r_{j+1}^{(s+z_j)} \geq r_j^{(s+z_j)}\) for \(0 \leq s \leq \alpha_{j+1}\). Computing \(T_{j-1}^{(s)}\) from \(T_{j-1}^{(s-1)}\) for \(1 \leq s \leq \alpha_j\), we denote \(b = \text{path}_j\) and \(q = \text{path}_j\) in Step 1 and Step 2 of Definition 3.19 by \(b_j^{(s)}\) and \(q_j^{(s)}\).

Since \((r_{j+1}, c_{j+1})\) is a corner cell in \(T_{j-1}^{(z_j)}\), we have \(r_{j+1}^{(0)} > r_j^{(0)}\). We prove that, for \(1 \leq s \leq \alpha_{j+1}\), we have that \(q_j^{(s)} > q_j^{(s+z_j)}\), which implies \(b_j^{(s)} > b_j^{(s+z_j)}\) and thus \(r_{j+1}^{(s)} > r_j^{(s+z_j)}\).

We prove \(q_j^{(s)} > q_j^{(s+z_j)}\) by induction on \(s\). First we check the case \(k = 1\). If \(r_{j+1}^{(0)} > r_j^{(z_j)}\), then it is obvious that \(q_j^{(1)} > q_j^{(z_j)}\). Otherwise if \(r_{j+1}^{(0)} = r_j^{(z_j)}\), we consider the following cases. \(q_j^{(z_j)}\) is the only element in \(A_{T_{j-1}^{(z_j)}}(r_{j+1}, c_{j+1})\). Let \(x = H_{T_{j-1}^{(z_j)}}(r_{j+1}, c_{j+1})\), \(y = \max(L_{T_{j-1}^{(z_j)}}(r_{j+1}, c_{j+1}))\) and \(y' = \max\{z \in L_{T_{j-1}^{(z_j)}}(r_{j+1}, c_{j+1}) \mid z \leq q_j^{(z_j)}\}\). See Figure 5 for illustration.

Case 1: If \(r_j^{(z_j)} = r_j^{(z_j)}\), then \(q_j^{(z_j)} = \).

Case 2: If \(r_j^{(z_j)} = r_j^{(z_j)}\), then \(q_j^{(z_j)} = \).

Since \(H_{T_{j-1}^{(0)}}(r_{j+1}, c_{j+1})\) is smaller or equal to \(y'\), we have that \(r_{j+1}^{(0)} \neq r_j^{(0)}\). Therefore \(q_j^{(1)}\) equals \(y\) when \(y > y'\) and \(q_j^{(z_j)}\) when \(y = y'\). In this case \(q_j^{(1)} = y = q_j^{(z_j)}\).

Now we have proved the base case \(s = 1\). Next, suppose it holds for some \(s \geq 1\) that \(q_j^{(s)} \geq q_j^{(s+z_j)}\) and \(r_j^{(s)} \geq r_j^{(s+z_j)}\). The statement is similar to the argument of the base case. If \(r_j^{(s)} > r_j^{(z_j)}\), it is
nonzero arm excess and
\[ c_j = \max(y^{(s+z_j)} \cup r^{(s+z_j)}). \]

Let \( x = H_{T_j}^{(s+z_j)}(r^{(s+z_j)}, c^{(s+z_j)}) \),
y = \max(z \in L^{(s+z_j)}(r^{(s+z_j)}, c^{(s+z_j)}) \ | \ z \leq q_j^{(s+z_j)}) \}. See
Figure 6 for illustration.

**Case (1):** If \( r_j^{(s+1+z_j)} = r_j^{(s+z_j)} \), then \( q_j^{(s+1+z_j)} = x \). If \( r_j^{(s+1)} = r_j^{(s)} \), then \( q_j^{(s+1)} = q_j^{(s+z_j)} \geq x \). If
\( r_j^{(s+1)} \neq r_j^{(s)} \), then \( q_j^{(s+1)} = \max(z \in L^{(s+z_j)}(r_j^{(s)}, c_j^{(s+z_j)}) \ | \ z \leq q_j^{(s+z_j)} \} \geq x \). So in either case
we have \( q_j^{(s+1)} \geq q_j^{(s+1+z_j)} \).

**Case (2):** If \( r_j^{(s+1+z_j)} \neq r_j^{(s+z_j)} \), then \( q_j^{(s+1+z_j)} = y' \). In this case we have \( H_{T_j}^{(s+z_j)}(r_j^{(s+z_j)} +
1, c_j^{(s+z_j)} - 1) \leq y' \leq q_j^{(s+z_j)} \). Since \( H_{T_j}^{(s)}(r_j^{(s+z_j)} + 1, c_j^{(s+z_j)} - 1) \) is smaller or equal to \( q_j^{(s+z_j)} \), we have
that \( r_j^{(s+1)} \neq r_j^{(s+1)} \). Therefore \( q_j^{(s+1)} = \max(z \in L^{(s+z_j)}(r_j^{(s)}, c_j^{(s+z_j)}) \ | \ z \leq q_j^{(s+z_j)} \} \leq q_j^{(s+1)} \}. By induction we have
\( q_j^{(s+z_j)} \leq q_j^{(s+1)} \), thus \( q_j^{(s+1)} \geq q_j^{(s+z_j)} \geq y' = q_j^{(s+1+z_j)} \). This completes the proof. \[ \square \]

**Lemma 3.26.** With the notations as above, let \( 0 \leq j \leq e - 1, 0 \leq s < \alpha_j + 1 \) and \( C_b([T_j^{(s)}], (r, c)]) =
[T_j^{(s+1)}, (r', c - 1)] \) for some \( r, c, r' \). Then in \( T_j^{(s+1)} \), column \( c - 1 \) is the rightmost column with nonzero arm excess and \( (r', c - 1) \) is the topmost cell in column \( c - 1 \) with nonzero arm excess.

**Proof.** In any path, consider the arm excess of its columns. Those with column index \( c \) such that
\( d(r_j, c_j) < c < c_j \) started with arm excess 0, then changed to arm excess 1 when the insertion path passed through that column, and immediately decreased to 0.

Thus the \( q_j^{(s)} \) that is being moved to cell \( (r', c - 1) \) is always at the rightmost column containing nonzero arm excess. When \( c - 1 > d(r_j, c_j) \), the arm excess of the column \( c - 1 \) is exactly 1, \( (r', c - 1) \)
is also the topmost cell containing an arm. For \( c - 1 = d(r_j, c_j) \), the path \( \text{path}_j \) has reached its destination. At that point, any column to the right of \( d(r_j, c_j) \) has 0 arm excess. It follows from Lemma 3.25 that the cell \((r_j^{(\alpha_j)}, c_j^{(\alpha_j)})\) is also the topmost cell containing an arm. \( \square \)

**Proposition 3.27.** The tableau \( T_j^{(s+1)} \) is a semistandard hook-valued tableau for all \( 0 \leq j \leq e - 1 \) and for all \( 0 \leq s < \alpha_j+1 \).

**Proof.** We only need to check that the \( q \) in Step 2 of Definition 3.19 is greater or equal to the hook entry and arm of the cell \( q \) is to be inserted into. When \( q \) is the only arm element, it is obvious that \( q \) is greater or equal to the hook entry.

The case when \( q \) is not the only arm element can only happen when we reach the destination column of the path. By the proof of Lemma 3.25, we have that for \( q_j^{(s)} \geq q_j^{(s+1)} \) for \( s \geq 1 \) and for \( j \) such that \( d(r_j, c_j) = d(r_j+1, c_j+1) \). Hence the statement follows by setting \( k = \alpha_j+1 \).

Before we define the “inverse” of the uncrowding map \( \mathcal{U} : \text{HVT}(\lambda) \to \sqcup_{\mu \geq \lambda} \text{SVT}(\mu) \times \hat{F}(\mu/\lambda) \), we need to restrict our domain to a subset \( K_\lambda \) of \( \sqcup_{\mu \geq \lambda} \text{SVT}(\mu) \times \hat{F}(\mu/\lambda) \), as the image of \( \mathcal{U} \) is not all of \( \sqcup_{\mu \geq \lambda} \text{SVT}(\mu) \times \hat{F}(\mu/\lambda) \). We define:

\[
K_\lambda(\mu) := \{(S, F) \in \text{SVT}(\mu) \times \hat{F}(\mu/\lambda) \mid \text{weight}(T_j^{(s)}) = \text{weight}(S), \forall 0 \leq j \leq e - 1, \forall 0 \leq s \leq \alpha_j+1\},
\]

\[
K_\lambda := \bigcup_{\mu \geq \lambda} K_\lambda(\mu).
\]

**Remark 3.28.** From the perspective of the uncrowding map, the set-valued tableau \( S \) in Example 3.20 cannot be obtained from a shape \((1,1)\) hook-valued tableau via the uncrowding map as explained in Remark 3.21. We say the cell \((1,2)\) in \( S \) practices social distancing. In this case,

\[
\begin{pmatrix}
3 \\ 2
\end{pmatrix}
\begin{pmatrix}
3 \\ 1 \\ 2
\end{pmatrix}
\begin{pmatrix}
1
\end{pmatrix}
\notin K_{(1,1)}.
\]

The \((S, F)\) in Example 3.24 is in \( K_{(3,2,1)}(5,3,2) \).

**Definition 3.29.** We can now define the crowding map \( \mathcal{C} \) for any partition \( \lambda \) as follows,

\[
\mathcal{C} : K_\lambda \to \text{HVT}(\lambda)
\]

\[
(S, F) \mapsto T_0^{(0)}.
\]

**Proposition 3.30.** The image of the uncrowding map \( \mathcal{U} : \text{HVT}(\lambda) \to \sqcup_{\mu \geq \lambda} \text{SVT}(\mu) \times \hat{F}(\mu/\lambda) \) is a subset of \( K_\lambda \). Moreover, we have \( \mathcal{C} \circ \mathcal{U} = \text{1}_{\text{HVT}(\lambda)} \).

**Proof.** We show that if \( \hat{h} = \mathcal{V}_h(h) \), where \( h \in \text{HVT} \), \( \mathcal{V}_h \) is as defined in Definition 3.2 and \( \hat{h} \) is obtained by moving some letter(s) from the cell \((r, c)\) to \((\hat{r}, \hat{c} + 1)\) (potentially adding a box), then \( C_\lambda([\hat{h}, (\hat{r}, \hat{c} + 1)]) = [h', (r', c)] \) satisfies \( [h', (r', c)] = [h, (r, c)] \).

We follow the notation used in Definitions 3.2 and 3.19. Thus \( a = \max(A_h(r,c)) \). We have that \( H_h(\hat{r}, \hat{c}) \leq a \). If cell \((r+1, c)\) is in \( h \), then \( H_h(r+1, c) > a \).

**Case (1):** \( \hat{r} \neq r \).

**Case (1A):** If cell \((\hat{r}, \hat{c} + 1)\) is not in \( h \), then \( h' \) is obtained by adding cell \((\hat{r}, \hat{c} + 1)\) and moving \( a \) from \( A_h(r,c) \) to \( H_h(\hat{r}, \hat{c} + 1) \). Under the action of \( C_\lambda \), by Step 1, \( b = a \) and \( r' = r \). \( C_\lambda \) appends \( a \) to \( A_h(\hat{r}, \hat{c} + 1) \) and removes cell \((\hat{r}, \hat{c} + 1)\), which recovers \( h \).

**Case (1B):** If cell \((\hat{r}, \hat{c} + 1)\) is in \( h \), then \( k \in L_r(\hat{r}, \hat{c} + 1) \) is the smallest number that is greater than or equal to \( a \) in column \( c + 1 \). \( h' \) is obtained by removing \( a \) from \( A_h(r,c) \), replacing \( k \) with \( a \), and attaching \( k \) to \( A_h(\hat{r}, \hat{c} + 1) \). Under the action of \( C_\lambda \), by Step 1, we can see that \( m = k, b = a \)
and \( r' = r \). By Step 3(b)i, \( q = b = a \), and \( a \) is appended to \( A_h^*(r, c) \) and \( q = a \) in \( L_h^*(\tilde{r}, c + 1) \) is replaced with \( m = k \). In the end, \( m \) is removed from \( A_h^*(\tilde{r}, c + 1) \). We recover \( h \).

Case (2): \( \tilde{r} = r \). Let \( \ell = \max(L_h^*(r, c)) \).

Case (2A): If cell \((r, c + 1)\) is not in \( h \), \( V_h \) adds cell \((r, c + 1)\), removes the part of \( L_h(r, c) \) that is greater than \( a \) to \( L_h(r, c + 1) \) and moves \( a \) from \( A_h(r, c) \) to \( H_h(r, c + 1) \). Under the action of \( C_h \), by Step 1, \( m = a \) and \( b = \ell \). Thus \( r' = r \). By Step 3(a)i, we move \( L_h^*(r, c + 1) \) into \( L_h^*(r, c) \) and we recover \( h \).

Case (2B): If cell \((r, c + 1)\) is in \( h \), \( \hat{h} \) is obtained by moving the part of \( L_h(r, c) \) that is greater than \( a \) to \( L_h(r, c + 1) \), moving \( a \) from \( A_h(r, c) \) to \( H_h(r, c + 1) \), and appending \( k \) to \( A_h(r, c + 1) \). Under the action of \( C_h \), by Step 1, \( m = a \) and \( b = \ell \). Then \( r' = r \) and \( q = a \). By Step 3(a)i, we move the set \( \{x \in L_h^*(r, c) \mid a < x \leq k \} \) from \( L_h^*(r, c + 1) \) into \( L_h^*(r, c) \), which is the set that was moved from cell \((r, c)\) by \( V_h \). Removing \( k \) from \( A_h(r, c + 1) \) and setting \( H_h(r, c + 1) = k \), we recover \( h \).

Now we have proven \( C_h([\tilde{h}, (\tilde{r}, c + 1)]) = [h', (r', c)] = [h, (r, c)] \). It follows that for any \((S, F) = U(h)\), we have that \( T_{j}^{(s)} \) is semistandard and has the same weight as \( S \) for all \( 0 \leq j \leq e - 1 \), for all \( 0 \leq s \leq \alpha_{j+1} \). Thus \( \text{image}(U) \subset K_\lambda \) and \( C \circ U = 1_{\text{HVT}(\lambda)} \). \( \square \)

Proposition 3.31. \( K_\lambda \) is a subset of the image of \( U : \text{HVT}(\lambda) \rightarrow \sqcup_{\mu \geq \lambda} \text{SVT}(\mu) \times \hat{\mathcal{F}}(\mu/\lambda) \). Moreover, \( U \circ C = 1_{K_\lambda} \).

Proof. Let \((S, F) \in K_\lambda \), then for all \( 0 \leq j < e \) and for all \( 0 \leq s < \alpha_{j+1} \), \( C_h(\tilde{T}_j^{(s)}, (r, c)) = T_j^{(s+1)}, (r', c - 1) \) for some \( r, c, r' \). We show that \( V_h(T_j^{(s+1)}) = T_j^{(s)} \) for all \( 0 \leq j < e \) and for all \( 0 \leq s < \alpha_{j+1} \). Following the notation in Definition 3.2, we first locate the rightmost column that contains nonzero arm excess, then determine the topmost cell in row \( \tilde{r} \) in that column with nonzero arm excess. We denote by \( a \) the largest arm element in that cell.

By Lemma 3.26, in \( T_j^{(s+1)} \), column \( c - 1 \) is the rightmost column with nonzero arm excess and \((r', c - 1)\) is the topmost cell in column \( c - 1 \) with nonzero arm excess.

Case (1): \( r' = r \). In this case either cell \((r + 1, c - 1)\) does not exist in \( T_j^{(s)} \), or \( H_j^{(s)}(r + 1, c - 1) > b \).

Case (1A): \( A_j^{(s)}(r, c) = \emptyset \). \( m = H_j^{(s)}(r, c) \) and \( b = \max(L_j^{(s)}(r, c)) \). Since \( r' = r \), \( q = m \), \( T_j^{(s+1)} \) is obtained by appending \( m \) to \( A_j^{(s)}(r, c - 1) \), moving \( L_j^{(s)}(r, c) \) into \( L_j^{(s)}(r, c - 1) \), and removing cell

![Figure 7](image_url)

**Figure 7.** Left: case (1A): \((\tilde{r}, c + 1)\) is not in \( h \). Right: case (1B): \((\tilde{r}, c + 1)\) is in \( h \).

![Figure 8](image_url)

**Figure 8.** Left: Case (1A): \((r, c + 1)\) is not in \( h \). Right: Case (1B): \((r, c + 1)\) is in \( h \).
(r, c) from $T_j^{(s)}$. Note that everything in $L_{T_j^{(s)}}(r, c)$ is greater than $m$ and everything in $L_{T_j^{(s)}}(r, c - 1)$ is smaller or equal to $m$.

For the $V_b$ action, we have $a = m$ and $b$ is the greatest letter in $L_{T_j^{(s+1)}}(r, c - 1)$. Since every letter in $T_j^{(s+1)}(r'', c)$ is smaller than $m$ for $r'' < r$, we have $\tilde{r} = r$. $V_b$ acts on $T_j^{(s+1)}$ by adding the cell $(r, c)$, setting the hook entry to be $m$, and moving $(m, b) \cap L_{T_j^{(s+1)}}(r, c - 1)$ to $L_{T_j^{(s+1)}}(r, c)$. Then we recover $T_j^{(s)}$.

**Figure 9.** Left: Case (1A): $A_{T_j^{(s)}}(r, c) = \emptyset$. Right: Case (1B): $A_{T_j^{(s)}}(r, c) \neq \emptyset$.

**Case (1B):** $A_{T_j^{(s)}}(r, c) \neq \emptyset$. $m$ is the only element in $A_{T_j^{(s)}}(r, c)$, $q = H_{T_j^{(s)}}(r, c)$ and $b = \max\{x \in L_{T_j^{(s)}} \mid x \leq m\}$. $T_j^{(s+1)}$ is obtained by appending $q$ to $A_{T_j^{(s)}}(r, c - 1)$, setting $H_{T_j^{(s)}}(r, c)$ to be $m$, deleting $A_{T_j^{(s)}}$, and moving $\{x \in L_{T_j^{(s)}} \mid r < x \leq m\}$ to $L_{T_j^{(s)}}(r, c - 1)$.

For the $V_b$ action, $a = q$ and $b$ is the greatest letter in $L_{T_j^{(s+1)}}(r, c - 1)$. Since every letter in $T_j^{(s+1)}(r'', c)$ is smaller than $q$ for $r'' < r$ and $m \geq q$, $\tilde{r} = r$. $V_b$ acts on $T_j^{(s+1)}$ by setting $H_{T_j^{(s+1)}}(r, c) = q$, $A_{T_j^{(s+1)}}(r, c) = m$, and moving $(q, b) \cap L_{T_j^{(s+1)}}(r, c - 1)$ to $L_{T_j^{(s+1)}}(r, c)$. We recover $T_j^{(s)}$.

**Case (2):** $r' \neq r$.

**Case (2A):** $A_{T_j^{(s)}}(r, c) = \emptyset$. Note that in this case, $C_b$ will move $m$ somewhere else and remove the cell $(r, c)$. Since weight$(T_j^{(s+1)}) = \text{weight}(T_j^{(s)})$, we must have that $L_{T_j^{(s)}}(r, c) = \emptyset$. So $b = q = m$. $T_j^{(s+1)}$ is obtained from $T_j^{(s)}$ by appending $m$ to $A_{T_j^{(s)}}(r', c - 1)$ and removing the cell $(r, c)$.

For the $V_b$ action, $a = m$. Since every letter in $T_j^{(s+1)}(r'', c)$ is smaller than $m$ for $r'' < r$, a new cell $(r, c)$ is added, $\tilde{r} = r$. $V_b$ acts on $T_j^{(s+1)}$ by moving $m$ to $H_{T_j^{(s+1)}}(r, c)$ and removing $m$ from $A_{T_j^{(s)}}(r, c)$.

**Figure 10.** Left: case (2A): $A_{T_j^{(s)}}(r, c) = \emptyset$. Right: case (2B): $A_{T_j^{(s)}}(r, c) \neq \emptyset$.

**Case (2B):** $A_{T_j^{(s)}}(r, c) \neq \emptyset$. $m$ is the only element in $A_{T_j^{(s)}}(r, c)$, $q = b = \max\{x \in L_{T_j^{(s)}}(r, c) \mid x \leq m\}$. $T_j^{(s+1)}$ is obtained by appending $b$ to $A_{T_j^{(s)}}(r', c - 1)$, replacing $b$ in $L_{T_j^{(s)}}(r, c)$ with $m$, and removing $m$ from $A_{T_j^{(s)}}(r, c)$. 
For the $\mathcal{V}_b$ action, $a = b$. Since every letter in $T^{(s+1)}_j(r'',c)$ is smaller than $b$ for $r'' < r$, $m$ is the smallest letter that is greater or equal to $b$ in column $c$. Hence $\tilde{r} = r$. $\mathcal{V}_b$ acts on $T^{(s+1)}_j$ by removing $b$ from $A^{(s+1)}_j(r',c-1)$, replacing $m$ in $L^{(s+1)}_j(r,c)$ with $b$, and attaching $m$ to $A^{(s+1)}_j(r,c)$. We recover $T^{(s)}_j$.

Therefore we have $\mathcal{V}_b(T^{(s+1)}_j) = T^{(s)}_j$ for all $0 \leq j \leq e - 1$, for all $0 \leq s \leq \alpha_j$, and $\mathcal{V}(T^{(0)}_{j+1}) = T^{(0)}_j$.

It follows that we also recover the recording tableau $F$. Thus $\mathcal{U}(T^{(0)}_e) = (S,F)$. □

**Corollary 3.32.** The uncrowding map $\mathcal{U}$ is a bijection between $\text{HVT}(\lambda)$ and $K_{\lambda}$ with inverse $\mathcal{C}$.

### 3.6. Alternative uncrowding on hook-valued tableaux

In Section 3.2, we defined an uncrowding map sending hook-valued tableaux to pairs of tableaux with one being set-valued and the other being column-flagged increasing. As hook-valued tableaux were introduced as a generalization of both set-valued tableaux and multiset-valued tableaux, it is natural to ask if there is an uncrowding map taking hook-valued tableaux to pairs of tableaux with one being multiset-valued. In this section we provide such a map.

**Definition 3.33.** The multiset uncrowding bumping $\hat{\mathcal{V}}_b$: $\text{HVT} \to \text{HVT}$ is defined by the following algorithm:

1. Initialize $T$ as the input.
2. If the leg excess of $T$ equals zero, return $T$.
3. Find the topmost row that contains a cell with nonzero leg excess. Within this column, find the cell with the largest value in its leg. (This is the rightmost cell with nonzero leg excess in the specified row.) Denote the row index and column index of this cell by $r$ and $c$, respectively. Denote the cell as $(r,c)$, its largest leg entry by $\ell$, and its rightmost arm entry by $a$.
4. Look at the row above $(r,c)$ (i.e. row $r+1$) and find the leftmost number that is strictly greater than $\ell$.
   - If no such number exists, attach an empty cell to the end of row $r+1$ and label the cell as $(r+1,\tilde{c})$, where $\tilde{c}$ is its column index. Let $k$ be the empty character.
   - If such a number exists, label the value as $k$ and the cell containing $k$ as $(r+1,\tilde{c})$ where $\tilde{c}$ is the cell’s column index.

   We now break into cases:
   (a) If $\tilde{c} \neq c$, then remove $\ell$ from $L_T(r,c)$, replace $k$ with $\ell$, and attach $k$ to the leg of $L_T(r+1,\tilde{c})$.
   (b) If $\tilde{c} = c$ then remove $[\ell,a] \cap A_T(r,c)$ from $A_T(r,c)$ where $[\ell,a] \cap A_T(r,c)$ is the multiset $\{z \in A_T(r,c) \mid \ell \leq z \leq a\}$. Remove $\ell$ from $L_T(r,c)$, insert $[\ell,a] \cap A_T(r,c)$ into $A_T(r+1,\tilde{c})$, replace the hook entry of $(r+1,\tilde{c})$ with $\ell$, and attach $k$ to $L_T(r+1,\tilde{c})$.
5. Output the resulting tableau.

**Definition 3.34.** The multiset uncrowding insertion $\hat{\mathcal{V}}_d$: $\text{HVT} \to \text{HVT}$ is defined as $\hat{\mathcal{V}}_d(T) = \hat{\mathcal{V}}_b^d(T)$, where the integer $d \geq 1$ is minimal such that $\text{shape}(\hat{\mathcal{V}}_b^d(T))/\text{shape}(\hat{\mathcal{V}}_b^{d-1}(T)) \neq \emptyset$ or $\hat{\mathcal{V}}_b^d(T) = \hat{\mathcal{V}}_b^{d-1}(T)$.

**Definition 3.35.** Let $T \in \text{HVT}(\lambda)$ with leg excess $\alpha$. The multiset uncrowding map $\hat{\mathcal{U}}: \text{HVT}(\lambda) \to \bigcup_{\mu \supset \lambda} \text{MVT}(\mu) \times \mathcal{F}(\mu/\lambda)$ is defined by the following algorithm:

1. Let $\hat{P}_0 = T$ and let $\hat{Q}_0$ be the flagged increasing tableau of shape $\lambda/\lambda$. 


(2) For $1 \leq i \leq \alpha$, let $\tilde{P}_{i+1} = \tilde{V}(\tilde{P}_i)$. Let $r$ be the index of the topmost row of $\tilde{P}_i$ containing a cell with nonzero leg excess and let $\tilde{r}$ be the row index of the cell $\text{shape}(\tilde{P}_{i+1})/\text{shape}(\tilde{P}_i)$. Then $\tilde{Q}_{i+1}$ is obtained from $\tilde{Q}_i$ by appending the cell $\text{shape}(\tilde{P}_{i+1})/\text{shape}(\tilde{P}_i)$ to $\tilde{Q}_i$ and filling this cell with $\tilde{r} - r$.

Define $\tilde{U}(T) = (\tilde{P}(T), \tilde{Q}(T)) := (\tilde{P}_\alpha, \tilde{Q}_\alpha)$.

**Example 3.36.** Let $T$ be the hook-valued tableau

\[
T = \begin{array}{c|c|c|c|c}
79 & 8 & & \\
233 & 78 & & \\
1 & 223 & 4 & \\
\end{array}
\]

Then, we obtain the following sequence of tableaux $\tilde{V}_b(T)$ for $0 \leq i \leq 2 = d$ when computing the first multiset uncrowding insertion:

\[
\begin{align*}
79 & \quad \rightarrow \quad 9 & \quad \rightarrow \quad 9 & \quad = \tilde{V}(T).
\end{align*}
\]

Continuing with the remaining multiset uncrowding insertions, we obtain the following sequences of tableaux for the multiset uncrowding map:

\[
\begin{align*}
79 & \quad \rightarrow \quad 9 & \quad \rightarrow \quad 9 & \quad \rightarrow \quad 9 & \quad = \tilde{P}(T),
\end{align*}
\]

\[
\begin{align*}
2 & \quad \rightarrow \quad 2 & \quad \rightarrow \quad 2 & \quad \rightarrow \quad 4 & \quad = \tilde{Q}(T).
\end{align*}
\]

**Proposition 3.37.** Let $T \in \text{HVT}$. Then $\tilde{U}(T)$ is well-defined.

**Proof.** The statement follows from a similar argument to the proofs found in Corollary 3.7 and Lemma 3.9.

Similar to the uncrowding map $U$, the multiset uncrowding map $\tilde{U}$ interwines with the corresponding crystal operators.

**Theorem 3.38.** Let $T \in \text{HVT}$.

(1) If $f_i(T) = 0$, then $f_i(\tilde{P}(T)) = 0$.
(2) If $e_i(T) = 0$, then $e_i(\tilde{P}(T)) = 0$.
(3) If $f_i(T) \neq 0$, we have $f_i(\tilde{P}(T)) = \tilde{P}(f_i(T))$ and $\tilde{Q}(T) = \tilde{Q}(f_i(T))$.
(4) If $e_i(T) \neq 0$, we have $e_i(\tilde{P}(T)) = \tilde{P}(e_i(T))$ and $\tilde{Q}(T) = \tilde{Q}(e_i(T))$.

**Proof.** The proof follows similarly to those found in Proposition 3.12, Lemma 3.13, and Theorem 3.14.
4. Applications

In this section, we provide the expansion of the canonical Grothendieck polynomials $G_\lambda(x; \alpha, \beta)$ in terms of the stable symmetric Grothendieck polynomials $G_\mu(x; \beta = -1)$ and in terms of the dual stable symmetric Grothendieck polynomials $g_\mu(x; \beta = 1)$ using techniques developed in [BM12]. We first review the basic definitions and Schur expansions of the two polynomials.

Recall from (1.1), that the stable symmetric Grothendieck polynomial is the generating function

$$G_\mu(x; -1) = \sum_{S \in SVT(\mu)} (-1)^{|S| - |\mu|} x^{\text{weight}(S)}.$$ 

Its Schur expansion can be obtained from the crystal structure on set-valued tableaux [MPS21]

$$G_\mu(x; -1) = \sum_{S \in SVT(\mu)} (-1)^{|S| - |\mu|} s^{\text{weight}(S)}.$$ 

**Definition 4.1.** The reading word $\text{word}(S) = w_1 w_2 \cdots w_n$ of a set-valued tableau $S \in SVT(\mu)$ is obtained by reading the elements in the rows of $S$ from the top row to the bottom row in the following way. In each row, first ignore the smallest element of each cell and read all remaining elements in descending order. Then read the smallest elements of each cell in ascending order.

**Example 4.2.** The reading word of $P(T)$ in Example 3.6 is $\text{word}(P(T)) = 8675423362111567$.

**Example 4.3.** The highest weight set-valued tableaux of shape $(2)$ are

$$
\begin{array}{c}
1 & 1 \\
1 & 2 \\
1 & 3 \\
1 & 4 \\
\end{array},
\begin{array}{c}
1 & 1 \\
2 & 1 \\
3 & 1 \\
4 & 1 \\
\end{array},
\begin{array}{c}
1 & 1 \\
1 & 2 \\
1 & 3 \\
\end{array},
\begin{array}{c}
1 & 1 \\
1 & 2 \\
2 & 1 \\
\end{array},
\begin{array}{c}
1 & 1 \\
2 & 1 \\
2 & 1 \\
\end{array},
\begin{array}{c}
1 & 1 \\
2 & 1 \\
2 & 1 \\
3 & 1 \\
\end{array},
\begin{array}{c}
1 & 1 \\
2 & 1 \\
2 & 1 \\
4 & 1 \\
\end{array},
\begin{array}{c}
1 & 1 \\
2 & 1 \\
2 & 1 \\
3 & 1 \\
4 & 1 \\
\end{array},
\end{array}
$$

which gives the Schur expansion

$$G_{(2)}(x; -1) = s_2 - s_21 + s_211 - s_2111 + \cdots.$$ 

The dual stable symmetric Grothendieck polynomials $g_\mu(x; 1)$ are dual to $G_\mu(x; -1)$ under the Hall inner product on the ring of symmetric functions.

**Definition 4.4.** A reverse plane partition of shape $\mu$ is a filling of the cells in the Ferrers diagram of $\mu$ with positive integers, such that the entries are weakly increasing in rows and columns. We denote the collection of all reverse plane partitions of shape $\mu$ by $RPP(\mu)$ and the set of all reverse plane partitions by $RPP$.

The evaluation $\text{ev}(R)$ of a reverse plane partition $R \in RPP$ is a composition $\alpha = (\alpha_i)_{i \geq 1}$, where $\alpha_i$ is the total number of columns in which $i$ appears. The reading word $\text{word}(R)$ is obtained by first circling the bottommost occurrence of each letter in each column, and then reading the circled letters row-by-row from top to bottom and left to right within each row.

**Example 4.5.** Consider the reverse plane partition

$$R = \begin{array}{c}
1 & 2 \\
1 & 1 \\
1 & 1 \\
3 & 1 \\
\end{array} \in RPP((3, 2)).$$

By circling the bottommost occurrence of each letter in each column, we obtain

$$R = \begin{array}{c}
1 & 2 \\
1 & 1 \\
1 & 1 \\
3 & 1 \\
\end{array},
\text{ev}(R) = (2, 1, 1), \text{ word}(R) = 2113.$$
Lam and Pylyavskyy [LP07] showed that the dual stable symmetric Grothendieck polynomials $g_{\mu}(x;1)$ are generating functions of reverse plane partitions of shape $\mu$

$$g_{\mu}(x;1) = \sum_{R \in \text{RPP}(\mu)} x^{\text{ev}(R)}.$$  

They also provided the Schur expansion of the dual stable symmetric Grothendieck polynomials [LP07, Theorem 9.8]

$$g_{\mu}(x;1) = \sum_{F} s_{\text{innershape}(F)},$$  

where the sum is over all flagged increasing tableaux whose outer shape is $\mu$.

**Example 4.6.** When $\mu = (\mu_1)$ is a partition with only one row, we have $g_{(\mu_1)}(x;1) = s_{(\mu_1)}$.

The flagged increasing tableaux of outer shape $(2,1,1)$ are

```
1 2 2
1
```

Hence $g_{211}(x;1) = s_{211} + 2s_{21} + s_2$.

According to [BM12], a symmetric function $f_\alpha$ over the ring $R$ is said to have a **tableaux Schur expansion** if there is a set of tableaux $T(\alpha)$ and a weight function $\text{wt}_\alpha : T(\alpha) \to R$ so that

$$f_\alpha = \sum_{T \in T(\alpha)} \text{wt}_\alpha(T)s_{\text{shape}(T)}.$$

Furthermore, any symmetric function with such a property has the following expansion in terms of $G_{\mu}(x;-1)$ and $g_{\mu}(x;1)$.

**Theorem 4.7.** [BM12, Theorem 3.5] Let $f_\alpha$ be a symmetric function with a tableaux Schur expansion

$$f_\alpha = \sum_{T \in T(\alpha)} \text{wt}_\alpha(T)s_{\text{shape}(T)}$$

for some $T(\alpha)$. Let $S(\alpha)$ and $R(\alpha)$ be defined as sets of set-valued tableaux and reverse plane partitions, respectively, by

$$S \in S(\alpha) \text{ if and only if } P(\text{word}(S)) \in T(\alpha), \text{ and}$$

$$R \in R(\alpha) \text{ if and only if } P(\text{word}(R)) \in T(\alpha),$$

where $P(w)$ is the RSK insertion tableau of the word $w$. We also extend $\text{wt}_\alpha$ to $S(\alpha)$ and $R(\alpha)$ by setting $\text{wt}_\alpha(X) := \text{wt}_\alpha(P(\text{word}(X)))$ for any $X \in S(\alpha)$ or $R(\alpha)$. Then we have

$$f_\alpha = \sum_{R \in R(\alpha)} \text{wt}_\alpha(R)G_{\text{shape}(R)}(x;-1), \text{ and}$$

$$f_\alpha = \sum_{S \in S(\alpha)} \text{wt}_\alpha(S)(-1)^{|S|-|\text{shape}(S)|}g_{\text{shape}(S)}(x;1).$$

**Proposition 4.8.** The canonical Grothendieck polynomials have a tableaux Schur expansion.

**Proof.** Recall the uncrowding map on set-valued tableaux of Definition 3.1

$$U_{\text{SVT}} : \text{SVT}(\mu) \longrightarrow \bigsqcup_{\nu \supseteq \mu} \text{SSYT}(\nu) \times F(\nu/\mu).$$

By Corollary 3.32, we have a bijection

$$U : \text{HVT}(\lambda) \to K_\lambda = \bigsqcup_{\mu \supseteq \lambda} K_\lambda(\mu).$$
Note that \( K_\lambda \subseteq \bigsqcup_{\mu \subseteq \lambda} \text{SVT}(\mu) \times \hat{F}(\mu/\lambda) \). Denote
\[
\phi_\lambda(S) = | \{ F \in \hat{F} \mid (S, F) \in K_\lambda \} |.
\]

Note that sometimes \( \phi_\lambda(S) = 0 \).

Given \( H \in \text{HVT}(\lambda) \), we have \( U(H) = (S, F) \in \text{SVT}(\mu) \times \hat{F}(\mu/\lambda) \) for some \( \mu \supseteq \lambda \) and \( |\mu| = |\lambda| + a(H) \). We can also obtain \( U_{\text{SVT}}(S) = (T, Q) \in \text{SSYT}(\nu) \times F(\nu/\mu) \) for some \( \nu \supseteq \mu \) and \( |\nu| = |H| \). The weights of \( H, S \) and \( T \) are the same. When \( H \) is highest weight, that is \( e_i(H) = 0 \) for all \( i \), then \( S \) and \( T \) are also of highest weight and weight \( (H) = \text{shape}(T) \). Denote by \( \text{HVT}_h(\lambda), \text{SVT}_h(\lambda), \text{SSYT}_h(\lambda) \) the subset of highest weight elements in \( \text{HVT}(\lambda), \text{SVT}(\lambda), \text{SSYT}(\lambda) \), respectively.

Applying [HS20, Theorem 4.6] and the above correspondence, we obtain
\[
G_\lambda(x; \alpha, \beta) = \sum_{H \in \text{HVT}_h(\lambda)} \alpha^{a(H)} \beta^{\ell(H)} s_{\text{weight}(H)} = \sum_{\mu \supseteq \lambda} \sum_{(S, F) \in K_\lambda(\mu)} \alpha^{|\mu| - |\lambda|} |\beta|s_{\text{weight}(S)}
\]
\[
= \sum_{\mu \supseteq \lambda} \sum_{S \in \text{SVT}_h(\mu)} \phi_\lambda(S) \alpha^{|\mu| - |\lambda|} |\beta|s_{\text{weight}(S)}
\]
\[
= \sum_{\mu \supseteq \lambda} \sum_{\nu \supseteq \mu} \sum_{T \in \text{SSYT}_h(\nu)} \sum_{Q \in F(\nu/\mu)} \phi_\lambda(U_{\text{SVT}}^{-1}(T, Q)) \alpha^{|\mu| - |\lambda|} |\beta|s_{\text{shape}(T)}
\]
\[
= \sum_{T \in \mathcal{T}(\lambda)} \text{wt}_\lambda(T) s_{\text{shape}(T)},
\]
where \( \mathcal{T}(\lambda) = \{ T \in \text{SSYT}_h(\nu) \mid \nu \supseteq \lambda \} \) and
\[
\text{wt}_\lambda(T) = \sum_{\mu : \lambda \subseteq \mu \subseteq \text{shape}(T)} \alpha^{|\mu| - |\lambda|} |\beta|s_{\text{shape}(T)} - |\mu| \sum_{Q \in F(\text{shape}(T)/\mu)} \phi_\lambda(U_{\text{SVT}}^{-1}(T, Q)).
\]

Note that Proposition 4.8 in particular implies that the canonical Grothendieck polynomials are Schur positive. This was known from [HS20], but here an explicit tableaux formula is given.

**Corollary 4.9.** The canonical Grothendieck polynomials have \( G_\mu(x; -1) \) and \( g_\mu(x; 1) \) expansions:
\[
G_\lambda(x; \alpha, \beta) = \sum_{R \in \text{R}(\lambda)} \text{wt}_\lambda(R) G_{\text{shape}(R)}(x; -1),
\]
\[
G_\lambda(x; \alpha, \beta) = \sum_{S \in \text{S}(\lambda)} \text{wt}_\lambda(S) (-1)^{|S| - |\text{shape}(S)|} g_{\text{shape}(S)}(x; 1).
\]

**Example 4.10.** We compute the first two terms in \( G_{(2)}(x; \alpha, \beta) = s_2 + \beta s_{21} + 2\alpha s_3 + 2\alpha\beta s_{31} + \cdots \).

The semistandard Young tableaux involved are
\[
\mathcal{T}((2)) = \left\{ \begin{array}{c}
1 & 1 & 1 & 1 \\
1 & 1 & 1 & 2 \\
1 & 1 & 1 & 1 & 2 \\
1 & 1 & 1 & 1 & 1 & \cdots
\end{array} \right\}.
\]
Labelling the tableaux $T_1, T_2, T_3, T_4, \ldots$, we have $\text{wt}_2(T_1) = 1, \text{wt}_2(T_2) = \beta, \text{wt}_2(T_3) = 2\alpha, \text{wt}_2(T_4) = 2\alpha\beta$. Next we compute the elements in $\mathbb{R}(2)$ and $S(2)$ that correspond to $T_1$ and $T_2$:

$$\{ R \in \mathbb{R}(2) \mid P(\text{word}(R)) = T_1 \} = \{ \begin{array}{c}
1 & 1 \\
1 & 1 \\
1 & 1 \\
1 & 1 \\
1 & 1 \\
1 & 1 \\
\ldots
\end{array} \}$$

$$\{ R \in \mathbb{R}(2) \mid P(\text{word}(R)) = T_2 \} = \{ \begin{array}{c}
2 & 1 & 1 \\
2 & 1 & 1 \\
2 & 1 & 1 \\
2 & 1 & 1 \\
2 & 1 & 1 \\
\ldots
\end{array} \}$$

$$\{ S \in S(2) \mid P(\text{word}(S)) = T_1 \} = \{ \begin{array}{c}
1 & 1 \\
\end{array} \}$$

$$\{ S \in S(2) \mid P(\text{word}(S)) = T_2 \} = \{ \begin{array}{c}
2 & 1 & 1 \\
\end{array} \}.$$

Applying the expansion formulas, we obtain

$$G_2(x; \alpha, \beta) = G_2(x; -1) + G_{2(1)}(x; -1) + G_{2(2)}(x; -1) + G_{2(11)}(x; -1) + \cdots$$

$$+ \beta(G_{2(1)}(x; -1) + G_{2(2)}(x; -1) + 2G_{2(11)}(x; -1) + \cdots) + \cdots$$

$$G_2(x; \alpha, \beta) = g_2(x; 1) + \beta(g_{21}(x; 1) - g_2(x; 1)) + \cdots.$$

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