SUMMARY   BlockSum, also known as KesanBlock in Japanese, is a Latin square filling type puzzle, such as Sudoku. In this paper, we prove that the decision problem whether a given instance of BlockSum has a solution or not is NP-complete.

key words: NP-completeness, combinatorial puzzle, Latin square, BlockSum

1. Introduction

In this paper, we show that BlockSum puzzle is computationally hard like many other combinatorial puzzles, e.g., Sudoku [1], Tetris [2], PuyoPuyo [3], [4], and classic Nintendo games [5]. Let \( n \) denote a natural number. An instance of BlockSum puzzle is given as an \( n \times n \) grid of empty cells such that the \( n^2 \) cells are partitioned into disjoint subsets and each subset is assigned an integer. Any subset of cells is connected, i.e., it consists of side-adjacent cells. We call a subset of connected cells a block. We call the integer assigned to a block the demand of the block. A player of BlockSum puzzle is asked to fill all the \( n^2 \) cells with integers in \( [n] = \{1, 2, \ldots, n\} \) so that the following conditions are satisfied.

Latin square condition: The integers assigned to the cells form an \( n \times n \) Latin square, i.e., in each row and in each column, every integer in \([n]\) appears exactly once.

Demand condition: In each block, the sum of the integers assigned to the cells equals to the demand of the block.

We show a BlockSum instance and its solution \((n = 4)\) in Fig. 1. In the figure, a block is indicated by a subset of cells surrounded by boldface lines, and its demand is indicated by a small digit. In Japan, BlockSum puzzle is often taken up in various media these days since Tetsuya Miyamoto, who is a successful private tutoring teacher, uses this puzzle for tutoring elementary school students [6], [7]. This motivates us to investigate whether or not BlockSum puzzle is essentially hard as other puzzles in computational sense. For related works on BlockSum puzzle, Haraguchi et al. recently proposed an algorithm that automatically produces BlockSum instances with various difficulty levels [8].

We show the NP-completeness of the decision problem version of BlockSum puzzle. Given a BlockSum instance, the decision problem asks to identify whether it has a solution or not. We refer to this decision problem simply as BlockSum. It is the following theorem that we intend to prove.

Theorem 1: BlockSum is NP-complete even if every block consists of at most 2 cells.

We give the proof of Theorem 1 by means of reduction from Monotone Not-All-Equal 3SAT, a well-known NP-complete problem [9]. Although the computational hardness of a recreational puzzle does not necessarily imply its amusingness, it is a common nature of widely accepted recreational puzzles. In the sense, our result might imply that BlockSum is not only useful for educational purpose but also potentially fascinating for puzzlers. Preparing terminologies, notations and lemmata in Sect. 2, we present the proof in Sect. 3. We give concluding remarks in Sect. 4.

2. Preliminaries

2.1 Latin Square

Suppose an \( n \times n \) grid of cells. We refer the cell in the \( i \)-th row and in the \( j \)-th column to \((i, j) \in [n] \times [n]\). We denote an assignment of integers in \([n]\) to the cells by a function \( \varphi : [n] \times [n] \rightarrow [n] \). The value \( \varphi(i, j) \) represents the integer assigned to cell \((i, j)\). We call \( \varphi \) an \( n \times n \) Latin square, or a Latin square of order \( n \), if, in each row and in each column, each integer in \([n]\) appears exactly once, i.e., \( \varphi(i, j) \neq \varphi(i, j') \) for any \( i, j, j' \in [n] \) \((j \neq j')\) and \( \varphi(i, j) \neq \varphi(i', j) \) for any \( i, i', j \in [n] \) \((i \neq i')\). We define the standard Latin square of
order $n$ as follows:

$$\varphi(i, j) = \begin{cases} 
  i - j + 1 & \text{if } i \geq j, \\
  i - j + 1 + n & \text{otherwise.}
\end{cases}$$

Figure 2 shows the standard Latin square of order 4.

Assume that $n = pq$ holds for two natural numbers $p$ and $q$. In the proof, we may partition the $n \times n$ grid of cells into $p^2$ equally sized square subgrids so that each subgrid has $q \times q$ cells. For any $a, b \in [p]$, we denote the subgrid from the cell $(q(a - 1) + 1, q(b - 1) + 1)$ to the cell $(q(a - 1) + q, q(b - 1) + q)$ by $S_{a,b}^{n,p}$. We define $S_{a,b}^{n,p}$ as a set of cells as follows:

$$S_{a,b}^{n,p} = \{(q(a - 1) + k, q(b - 1) + \ell) \mid q = n/p, k, \ell \in [q]\}.$$  

(1)

Let us introduce a systematic way to construct a Latin square satisfying a certain condition. Let us denote a $p \times p$ Latin square by $\pi$. Let us denote a set of $p^2$ Latin squares of order $q$ by $\Psi = [\psi_{a,b} : [q] \times [q] \rightarrow [q] \mid (a, b) \in [p] \times [p]]$. We define $\varphi : [n] \times [n] \rightarrow [n]$ as the integer assignment that is obtained by pasting $\psi_{a,b} \in \Psi$ to each subgrid $S_{a,b}^{n,p}$ and by adding $q(\pi(a, b) - 1)$ to the integers, i.e., for any $a, b \in [p]$ and any $k, \ell \in [q]$,

$$\varphi(q(a - 1) + k, q(b - 1) + \ell) = \psi_{a,b}(k, \ell) + q(\pi(a, b) - 1).$$  

(2)

**Proposition 2:** The $\varphi$ given by (2) is an $n \times n$ Latin square.

**Proof:** We show that $\varphi$ satisfies the Latin square condition. Clearly, the integers that $\varphi$ assigns to the $n \times n$ grid are in $[n]$. Let us take any two cells $(i, j)$ and $(i', j')$ that are in the same $i$-th row. If they are in the same subgrid $S_{a,b}^{n,p}$, they are assigned different values by $\varphi$ since each subgrid is pasted a Latin square and all the integers in the subgrid are added the equivalent value, that is $q(\pi(a, b) - 1)$. Otherwise, i.e., if $(i, j)$ and $(i', j')$ are in different subgrids $S_{a,b}$ and $S_{a,b'}$, respectively, the two cells are assigned different values since $\varphi$ assigns the integers of disjoint ranges to these subgrids due to $\pi(a, b) \neq \pi(a, b')$. Analogously, any two cells in the same column are assigned different values. $\square$

**Lemma 3:** For any natural number $z$, there is a $2^z \times 2^z$ Latin square $\varphi : [2^z] \times [2^z] \rightarrow [2^z]$ such that

$$\varphi(i, j) = \begin{cases} 
  1 & \text{if } i = j, \\
  j & \text{if } i \neq j \text{ and } i = 1, \\
  i & \text{if } i \neq j \text{ and } j = 1.
\end{cases}$$  

(3)

**Proof:** Prove by induction. When $z = 1$, we have such a Latin square, as shown in Fig. 3. For a general $z \geq 2$, assume that we have a $2^{z-1} \times 2^{z-1}$ Latin square that satisfies (3). Let us denote it by $\pi_{z-1}$. Then it is easy to see that a Latin square constructed by (2) with $p = 2$, $q = 2^{z-1}$, $\pi = \pi_z$, and $\psi_{a,b} = \pi_{z-1}$ for any $(a, b) \in [2] \times [2]$ satisfies (3). This and Proposition 2 show the lemma. (For example, see Fig. 3 for $z = 2$ and 3. We can see that shaded cells surely satisfy (3).) $\square$

### 2.2 BlockSum

Two cells $(i, j)$ and $(i', j')$ are *adjacent* if $|i - i'| + |j - j'| = 1$. The adjacency defines the connectivity of cells. A block is a set of connected cells. Let us denote a block by $B \subseteq [n] \times [n]$. We call $B$ a *k-block* if $|B| = k$. In particular, we call $B$ an *$(r \times c)$-block* if it is an $(rc)$-block such that the cells form an $r \times c$ rectangle. Let us denote $\varphi : [n] \times [n] \rightarrow [n]$. When a block $B$ is given by $B = \{(i_1, j_1), \ldots, (i_k, j_k)\}$, we denote by $\varphi(B)$ the sequence $(\varphi(i_1, j_1), \ldots, \varphi(i_k, j_k))$ for convenience.

We represent a BlockSum instance by $I_{BS} = (B, \sigma)$, where $B$ denotes a partition of the $n^2$ cells into blocks and $\sigma$ denotes a function $\sigma : B \rightarrow \{n^2(n + 1)/2\}$. For any block $B \in B$, the value $\sigma(B)$ denotes the demand of $B$. We call $\sigma$ a *demand function*. Then the decision problem BlockSum is formally defined as follows.

**BlockSum**

**Instance:** A BlockSum instance $I_{BS} = (B, \sigma)$, where $B$ is a partition of the $n^2$ cells into blocks and $\sigma$ is a demand function $\sigma : B \rightarrow \{n^2(n + 1)/2\}$.

**Question:** Is there an assignment $\varphi : [n] \times [n] \rightarrow [n]$ such that $\varphi$ is a Latin square (the Latin square condition) and $\sum_{(i,j) \in B} \varphi(i, j) = \sigma(B)$ holds for any $B \in B$ (the Demand condition)?

When the answer is “yes,” we say that the instance $I_{BS}$ is
solvable and that the assignment $\varphi$ satisfying the two conditions is a solution of $I_{BS}$.

Here, we show a technique of transferring a BlockSum instance by exchanging demands that will be used in the proof of NP-completeness. Suppose that, in $I_{BS} = (\mathcal{B}, \sigma)$, four $2 \times 1$ blocks $B_1, B_2, B_3, B_4$ are in $\mathcal{B}$ such that they form

$$B_1 = \{(r, c), (r + 1, c), B_2 = \{(r, c'), (r + 1, c'), B_3 = \{(r', c), (r + 1, c), B_4 = \{(r', c'), (r' + 1, c'), \}$$

and $\sigma(B_1) = \sigma(B_4) = s$ and $\sigma(B_2) = \sigma(B_3) = t$. In Fig. 4, we illustrate how $B_1, \ldots, B_4$ are located in $I_{BS}$. Let us construct another instance by exchanging the demands between $B_1$ and $B_2$ and the demands between $B_3$ and $B_4$. That is, denoting the constructed instance by $I'_{BS} = (\mathcal{B}, \sigma')$, we define the demand function $\sigma'$ as follows:

$$\sigma'(B) = \begin{cases} t & \text{if } B = B_1 \text{ or } B_4, \\ s & \text{if } B = B_2 \text{ or } B_3, \\ \sigma(B) & \text{otherwise.} \end{cases}$$

Assume that $I_{BS}$ has a solution $\varphi : [n] \times [n] \to [n]$ such that

$$\varphi(B_1) = \varphi(B_4) = (s_1, s_2), \varphi(B_2) = \varphi(B_3) = (t_1, t_2),$$

and $s_1 + s_2 = s, t_1 + t_2 = t$. Clearly, the following $\varphi' : [n] \times [n] \to [n]$ is a solution of $I'_{BS}$:

$$\begin{align*}
\varphi'(B_1) &= \varphi'(B_4) = (s_1, s_2), \\
\varphi'(B_2) &= \varphi'(B_3) = (t_1, t_2), \\
\varphi'(i, j) &= \varphi(i, j) \text{ for any } (i, j) \notin B_1 \cup \cdots \cup B_4. 
\end{align*}$$

Lemma 4: If $\varphi$ is a solution of $I_{BS}$ that satisfies (4), the assignment $\varphi'$ given by (5) is a solution of $I'_{BS}$.

2.3 Monotone Not-All-Equal 3SAT

Let $V = \{v_1, v_2, \ldots, v_N\}$ denote a set of $N$ Boolean variables. For a variable $v \in V$, $v$ is called the positive literal and $\bar{v}$ is called the negative literal. A clause is a set of literals over $V$. A truth assignment for $V$ is denoted by $\tau : V \to \{\text{True, False}\}$. Given a true assignment $\tau$, if $\tau(v)$ = True, then the positive literal $v$ is true and the negative literal $\bar{v}$ is false. If $\tau(v) = \text{False}$, then $v$ is false and $\bar{v}$ is true. A clause is called not-all-equal under $\tau$ if the clause has at least one literal that is true under $\tau$ and at least one literal that is false under $\tau$. The decision problem Not-All-Equal SAT is defined as follows.

**Not-All-Equal SAT**

**Instance:** A Not-All-Equal SAT instance $I_{SAT} = (V, C)$, where $V$ is a set $V = \{v_1, v_2, \ldots, v_N\}$ of $N$ Boolean variables, and $C$ is a collection $C = \{C_1, C_2, \ldots, C_M\}$ of $M$ clauses over $V$.

**Question:** Is there a truth assignment $\tau : V \to \{\text{True, False}\}$ such that each $C_b \in C$ is not-all-equal under $\tau$?

We abbreviate Not-All-Equal SAT into NAE-SAT. When the answer is “yes,” we call $I_{SAT}$ NAE-satisfiable. The problem NAE-3SAT is the special case of NAE-SAT such that $|C_b| = 3$ for any $C_b \in C$. The NAE-3SAT is known to be NP-complete [10]. Further, the problem Monotone NAE-3SAT is the special case of NAE-3SAT such that all literals are positive. The Monotone NAE-3SAT is still NP-complete [9].

**Lemma 5:** NAE-SAT is NP-complete even if we restrict an instance $I_{SAT} = (V, C)$ so that

(i) any variable in $V$ appears as 2 positive literals and as 1 negative literal in $C$,

(ii) any clause in $C$ has either 2 or 3 literals, and

(iii) any clause of 2 literals has one positive literal and one negative literal, and any clause of 3 literals has only positive literals.

Proof: We construct a polynomial transformation from a Monotone NAE-3SAT instance $I_{SAT} = (W, D)$ to a NAE-SAT instance $I_{SAT} = (V, C)$ so that it satisfies (i), (ii) and (iii). Let us denote $N = |W|$ and $M = |D|$. Suppose that, in $D$, a Boolean variable $w_a \in W$ appears as a positive literal $k_a$ times. We take a set of $k_a$ Boolean variables, denoted by $X_a = \{x_{a,1}, x_{a,2}, \ldots, x_{a,k_a}\}$. Taking the set $X_a$ for each Boolean variable $w_a \in W$, we define $V = X_1 \cup \cdots \cup X_N$. We have $|V| = \sum_{a=1}^{N} k_a = 3M$.

We generate a set of clauses over $V$ as follows. For each Boolean variable $w_a \in W$, we replace all $k_a$ positive literals $w_a$-s in $D$ with $k_a$ positive literals $x_{a,1}, x_{a,2}, \ldots, x_{a,k_a}$, respectively. We denote the set of clauses generated in this way by $C^{(3)}$. We have $|C^{(3)}| = |D| = M$. Each clause in $C^{(3)}$ has exactly 3 literals since each clause in $D$ does so. For any variable in $V$, its positive literal appears in $C^{(3)}$ exactly once. We do not have any other literals in $C^{(3)}$.

Next, we take another set of clauses over $V = X_1 \cup \cdots \cup X_N$. The set is denoted by $C^{(2)} = C^{(2)}_1 \cup \cdots \cup C^{(2)}_N$, where each $C^{(2)}_b$ is a set of clauses over $X_a$ for the variable $w_a \in W$ as follows:

$$C^{(2)}_b = (\bigcup_{b=1}^{k_a} \{[x_{a,b}, \bar{x}_{a,b+1}]\}) \cup \{[x_{a,k_a}, \bar{x}_{a,1}]\}.$$ (6)

We have $|C^{(2)}| = \sum_{a=1}^{N} k_a = 3M$. Clearly, each clause in $C^{(2)}$ has exactly 2 literals. For any variable in $X_a$, its positive and
negative literals appear in $C_2^{(2)}$ exactly once respectively.

The NAE-SAT instance $I_{SAT} = (V, C)$ with $V = X_1 \cup \cdots \cup X_N$ and $C = C^{(2)} \cup C^{(3)}$ satisfies (i), (ii) and (iii) from the discussion so far. We have only to show that the above procedure is a polynomial transformation. Clearly, the construction can be done in a polynomial time in $N$ and $M$. We claim that $J_{SAT}$ is equivalent to $I_{SAT}$, i.e., $J_{SAT}$ is NAE-satisfiable if and only if $I_{SAT}$ is NAE-satisfiable. To show the necessity, suppose a truth assignment $\tau : W \rightarrow \{\text{True}, \text{False}\}$ such that each clause in $D$ is not-all-equal. From $\tau$, we construct a truth assignment $\omega : V \rightarrow \{\text{True}, \text{False}\}$ as follows; for each variable $x$ in $X_a$, we let $\omega(x) = \tau(w_a)$. Thus all the variables in the same $X_a$ are assigned the same truth value. From (6), we see that any clause in $C^{(2)} = C^{(2)}_1 \cup \cdots \cup C^{(2)}_N$ is not-all-equal under $\omega$. Recall that any clause $C \in C^{(3)}$ has been generated from a certain clause $D \in D$, by replacing literals over $W$ with ones over $V$. From the assumption, there is a positive literal $u \in D$ that is true under $\tau$. If $u$ is a positive literal of a variable $w_b \in W$, then $\tau(w_b)$ should be True. In $C$, $u$ is replaced with a certain literal $x_{a,b}$. Since $\omega(x_{a,b}) = \tau(w_b) = \text{True}$, $C$ has a literal that is true under $\omega$. Similarly, the clause $D$ also has a literal that is false under $\tau$, which guarantees that $C$ has a literal that is false under $\omega$. Hence, $C$ is not-all-equal under $\omega$. We are done with the necessity. The sufficiency can be shown in an analogous way. \( \square \)

3. Proof of NP-Completeness of BlockSUM

Now we are ready to prove Theorem 1, the NP-completeness of BlockSUM. It is easy to see that BlockSUM is in NP; given a BlockSUM instance and an integer assignment $\varphi : [n] \times [n] \rightarrow [n]$, we can verify whether or not $\varphi$ is a solution of the instance in $O(n^2)$ time.

We reduce NAE-SAT to BlockSUM. In the reduction, we establish a polynomial transformation from any NAE-SAT instance satisfying (i), (ii) and (iii) of Lemma 5 to a BlockSUM instance such that each block consists of at most 2 cells. Let us denote an NAE-SAT instance satisfying (i), (ii) and (iii) of Lemma 5 by $I_{SAT} = (V, C)$, where $V = \{v_1, v_2, \ldots, v_N\}$ is the set of $N$ Boolean variables and $C = \{C_1, C_2, \ldots, C_M\}$ is the set of $M$ clauses over $V$.

We construct a BlockSUM instance $I_{BS} = (B, \sigma)$ on a $(10MN_0) \times (10MN_0)$ grid of cells, where $N_0 = 2^{[\log_2(N+1)]}$. Since we have

$$N_0 = 2^{[\log_2(N+1)]} \leq 2^{\log_2(N+1)+1} = 2(N+1),$$

the number of cells in the grid is at most $(20M(N+1))^2$, a polynomial in $M$ and $N$. Let us denote $n = 10MN_0$. The $n \times n$ grid is partitioned into $N_0^2$ equally sized square subgrids so that each subgrid has $10M \times 10M$ cells. Recall that the notation $S_{a,b}^{n,n_0}$ represents a subgrid defined by (1). In the sequel, we let $S_{a,b}^{n,n_0}$ represent $S_{a,b}^{n,n_0}$ for convenience if no confusion arises. Any block in $B$ is contained in one subgrid, or equivalently, no block crosses more than one subgrid. Each subgrid consists of 1-blocks and 2-blocks. Let the caligraphic

![Fig 5 Overview of the n \times n grid of the BlockSum instance to be constructed (n = 10MN_0).](image)

$S_{a,b}$ denote the partition of a subgrid $S_{a,b}$ into blocks. We take the partition $B$ as the union of the partitions of the subgrids:

$$B = \bigcup_{(a,b) \in [N_0] \times [N_0]} S_{a,b}.$$  

For the demand function $\sigma$, we first set $\sigma$ so that the instance $I_{BS} = (B, \sigma)$ surely has a solution, regardless of the satisfiability of $I_{SAT}$. Then we will exchange the demands of some blocks by means of Lemma 4 so that the resulting instance has a solution if and only if $I_{SAT}$ is NAE-satisfiable. To guarantee that $I_{BS}$ has a solution, we construct an $n \times n$ Latin square somehow, denoted by $\varphi$ and set the demand $\sigma(B)$ of each block $B \in B$ as follows:

$$\sigma(B) = \sum_{(i,j) \in B} \varphi(i,j).$$  

Clearly, the BlockSUM instance $(B, \sigma)$ has $\varphi$ as a solution. For $\varphi$, we take such an $n \times n$ Latin square that can be given by the pasting and adding procedure of Proposition 2; since $N_0 = 2^{[\log_2(N+1)]}$, there is an $N_0 \times N_0$ Latin square that satisfies (3) of Lemma 3, denoted by $\pi$. We can construct such $\pi$ in $O(N^2)$ time. Then the $n \times n$ Latin square $\varphi$ is constructed by (2), with $\pi$ and $10M \times 10M$ Latin squares $\psi_{a,b}$ that are appropriately chosen for each subgrid $S_{a,b}$. The point is that we can choose the Latin square $\psi_{a,b}$ for the subgrid $S_{a,b}$ independently from other subgrids.

Now we explain how we define the partition $S_{a,b}$ and the Latin square $\psi_{a,b}$ for each subgrid $S_{a,b}$ that determines the demands of the blocks along with $\pi$. We illustrate the roles of the subgrids in Fig. 5. We use $S_{1,1}$ as the satisfaction testing component. In other words, we embed the blocks into $S_{1,1}$ that correspond to the $M$ clauses. For any $a \in [N]$, we use $S_{a,1+1}$, $S_{a+1,1}$ and $S_{a+1,a+1}$ as the truth assignment components for the variable $v_a$. These subgrids are
contained in our grid properly since we have
\[ N + 1 = 2^{\log_2(N+1)} \leq 2^{\log_2(2N)} = 2N. \]

We use the other \( N_0 - 1 \) subgrids for garbage collection.

(1) Satisfaction Testing Components

We use the \( 10M \times 10M \) subgrid \( S_{1,1} \) as the satisfaction testing component. Let us partition \( S_{1,1} \) further into \( M^2 \) subgrids so that each subgrid has \( 10 \times 10 \) cells. We focus on the \( M \) subgrids on the diagonal. We denote the \( M \) subgrids by \( U_1, U_2, \ldots, U_M \). Formally, each \( U_b \) is defined as \( U_b = S_{b,b}^{10M} (b \in [M]) \). We illustrate how \( U_1, U_2, \ldots, U_M \) are located in Fig. 6. A subgrid \( U_b \) corresponds to the clause \( C_b \in C \). When \( C_b \) has 2 literals (resp., 3 literals), we take the cells of \( U_b \) and the demands of the blocks as shown in Fig. 7 (resp., Fig. 8). The cell of a 1-block should be filled with its demand, and we cannot assign the demand value to any other cell in the same row or in the same column. For example, in Fig. 7, there are 1-blocks whose demands are from 5 to 10 in the 1st to 4th rows and in the 1st to 4th columns. Thus we cannot assign 5 to 10 to the upper-left \( 4 \times 4 \) subgrid, and thus need to assign 1 to 4 there. We do not take up the gray 1-blocks any more. We associate each literal in \( C_b \) with a certain cell in \( U_b \), as shown in the figures. When \( |C_b| = 2 \), the clause has one positive literal and one negative literal from Lemma 5 (iii). We denote the cell for the positive (resp., negative) literal by \( u_b^{(2)} \) (resp., \( u_b^{(3)} \)). When \( |C_b| = 3 \), the clause has three positive literals, and we denote the cells for these literals by \( u_b^{(1)}, u_b^{(2)}, u_b^{(3)} \) and \( u_b^{(3)} \).

When \( |C_b| = 2 \), the 2-blocks in the subgrid \( U_b \) admits only 2 configurations of integers, as shown in Fig. 9. When \( |C_b| = 3 \), the 3-blocks in the subgrid \( U_b \) admits 6 \( \times 2^3 = 96 \) configurations, as shown in Fig. 10; the cells from the 1st row to the 6th row admit only 6 configurations. We can exchange the integers assigned to two out of the four \( (2 \times 1) \) blocks in the 7th and 8th rows. For example, the assignment of \((1, 1)\) and \((2, 4)\) to the two \((2 \times 1)\) blocks can be flipped into \((4, 1)\) and \((1, 4)\). This is also true for the four \((2 \times 1)\) blocks in the 9th and 10th rows. Let us emphasize that, whether \( |C_b| = 2 \) or 3, the cells for the literals are assigned either 1 or 2, and they are not-all-equal.

We let every cell out of \( U_1, \ldots, U_M \) form a 1-block. In any feasible configuration of integers to \( U_1, \ldots, U_M \), the integers from 1 to 10 are assigned in all rows and in all columns of \( S_{1,1} \). We readily see that we can assign integers from 11 to 10M to the cells out of \( U_1, \ldots, U_M \) so that the Latin square condition is satisfied. We use the assigned integers as the demands of the 1-blocks. Then, the configuration of the integers to these cells is unique.

(2) Truth Assignment Components

For each \( a \in [N] \), we use the subgrids \( S_{a,1}, S_{a+1,1} \) and \( S_{a+a,1} \) as the truth assignment components for the variable \( v_a \in V \). From Lemma 5 (i), \( v_a \) appears as 2 positive literals and as 1 negative literal in \( C \). These 3 literals are associated with some 3 cells in \( S_{1,1} \). We denote the 2 cells for the 2 positive literals by \( u_a^{(1)} \) and \( u_a^{(2)} \), and the cell for the negative literal by \( u_{a,3} \). As shown in Figs. 7 and 8, these cells belong to \((2 \times 1)\) blocks. We denote the three \((2 \times 1)\) blocks to which
blocks in the subgrid $S_{a+1}$, the demands of all the blocks in $S_{a+1}$ are determined by (7). In particular, the demands of $D_{a,1}$, $D_{a,2}$ and $D_{a,3}$ becomes:

$$\sigma(D_{a,1}) = \sigma(D_{a,2}) = \sigma(D_{a,3}) = (10Ma + 1) + (10Ma + 2) = 20Ma + 3. \quad (8)$$

Note that the configuration of integers to the subgrid $S_{a+1}$ is unique; the 1-blocks need to be assigned their demands. The 2-block $D_{a,2}$ ($t = 1, 2, 3$) must be assigned $(10Ma + 1, 10Ma + 2)$ since, in the $r_{t}$-th row (resp., in the $(r_{t} + 1)$-th row), all the integers from 1 to $n$ except $10Ma + 1$ (resp., $10Ma + 2$) need to be assigned to other cells other than $D_{a,2}$.

We set the partition and the demands of the subgrid $S_{a+1,1}$ in the similar way. The subgrid $S_{a+1,1}$ has three (2 x 1)-blocks, and all the other blocks are 1-blocks. We denote the three (2 x 1)-blocks by $D'_{a,1}$, $D'_{a,2}$ and $D'_{a,3}$, where we define

$$D'_{a,1} = \{(10Ma + 1, c_1), (10Ma + 2, c_1)\},$$
$$D'_{a,2} = \{(10Ma + 3, c_2), (10Ma + 4, c_2)\},$$
$$D'_{a,3} = \{(10Ma + 5, c_3), (10Ma + 6, c_3)\}. $$

Thus, these 3 blocks are in the same columns as $A_{a,1}$, $A_{a,2}$ and $A_{a,3}$, respectively. We construct $\psi_{a+1,1}$ by permuting the columns of the standard Latin square of order 10 so that, in the $c_1$-th, $c_2$-th and $c_3$-th columns, 1 is in the $r_1$-th, $r_2$-th and $r_3$-th rows respectively. One can readily see that this permutation is always feasible. With $\pi$ and $\psi_{a+1,1}$, the demands of all the blocks in $S_{a+1,1}$ are determined by (7). In particular, the demands of $D'_{a,1}$, $D'_{a,2}$ and $D'_{a,3}$ becomes:

$$\sigma(D'_{a,1}) = \sigma(D'_{a,2}) = \sigma(D'_{a,3}) = 20Ma + 3. \quad (9)$$

Note that the configuration of integers to the subgrid $S_{a+1,1}$ is unique.
Thus the configuration of integers to $S_{a,b}$ is unique.

Finally, we transfer $I'_{BS}$ to another instance $I'_{BS} = (B, \sigma')$ by exchanging the demands of some blocks by means of Lemma 4, while we do not change the partition $B$. Again, see Fig. 11. For each Boolean variable $v_a \in V$ and $t = 1, 2, 3$, $D_{a,t}$ in $S_{a,t+1}$ is in the same rows as $A_{a,t}$ in $S_{a,t}$ and $D'_{a,t}$ in $S_{a+1,t+1}$ is in the same columns as $A_{a,t}$. The $E_{a,t}$ in $S_{a+1,t+1}$ is in the same columns as $D_{a,t}$ and in the same rows as $D'_{a,t}$. Also we have $\sigma(A_{a,t}) = \sigma(E_{a,t}) = 5$ and $\sigma(D_{a,t}) = \sigma(D'_{a,t}) = 20Ma + 3$, where the latter is from (8) and (9). Then we define the demand function $\sigma' : B \rightarrow [n^2(n + 1)/2]$ as follows; for any $a = 1, 2, \ldots, N$ and $t = 1, 2, 3$, we define:

$\sigma'(A_{a,t}) = \sigma'(E_{a,t}) = 20Ma + 3$,
\[ \sigma'(D_{a,t}) = \sigma'(D'_{a,t}) = 5, \]
and $\sigma'(B) = \sigma(B)$ for any other block $B$ in $B$.

We have finished explaining the transformation from $I_{SAT}$ to $I'_{BS}$. Note that every block in $I'_{BS}$ has at most 2 cells. We can readily see that the transformation time is bounded by a polynomial in $M$ and $N$. Next we show that $I'_{BS}$ has a solution if and only if $I_{SAT}$ is NAE-satisfiable.

Suppose that $I_{SAT}$ is NAE-satisfiable. Let $\tau : V \rightarrow \{\text{TRUE}, \text{FALSE}\}$ denote a truth assignment such that each clause $C_b$ in $C$ is not-all-equall. From $\tau$, we first construct a solution $\varphi$ of $I_{BS}$ and then transform it into a solution of $I'_{BS}$ by means of Lemma 4. For $I_{BS}$, the configuration to the cells out of $S_{1,1}, S_{2,2}, \ldots, S_{N+1,N+1}$ is unique. Also, for $S_{1,1}$, the configuration to the cells out of $U_{1,1}, \ldots, U_{N+1,N+1}$ is unique, and for $S_{a+1,t+1}$ ($\forall a \in [N]$), the configuration to the cells out of the upper-left $10 \times 10$ subgrid is unique. In $U_b$ ($b \in [M]$) that corresponds to the clause $C_b$, there are 2 or 3 cells with which the literals in $C_b$ are associated. Assign 2 (resp., 1) to the cell if the literal is true (resp., false) under $\tau$. Since $C_b$ is not-all-equall, the assigned integers are not-all-equall. Therefore, when $|C_b| = 2$ (resp., 3), we can employ one of the configurations shown in Fig. 9 (resp., Fig. 10). In the upper-left $10 \times 10$ subgrid of $S_{a+1,t+1}$ ($\forall a \in [N]$), assign (2, 2, 1) (resp., (1, 1, 2)) to the cells $w_{a,1}, w_{a,2}$ and $w_{a,3}$ if $\tau(v_a) = \text{TRUE}$ (resp., FALSE), which leads to one of the configurations shown in Fig. 13. We easily see that, for $t = 1, 2, 3$, $\varphi(A_{a,t}) = \varphi(E_{a,t})$ and $\varphi(D_{a,t}) = \varphi(D'_{a,t})$ hold. Therefore, this $\varphi$ can be transformed into a solution of $\varphi'$ by means of Lemma 4.

Conversely, suppose that $I'_{BS}$ is solvable. Let $\varphi'$ denote any solution of $I'_{BS}$. For any $a = 1, 2, \ldots, N$, let us see the $2 \times 1$-blocks $D_{a,1}', D_{a,2}'$ and $D_{a,3}'$ in the subgrid $S_{a+1,1}$ whose demands by $\sigma'$ are 5. Let $w_{a,1}'$ and $w_{a,2}'$ denote the lower cells of the $D_{a,1}'$ and $D_{a,2}'$ respectively and $w_{a,3}'$ denote the upper cell of $D_{a,3}'$.

**Lemma 6:** We have $\varphi'(w_{a,t}') \in \{1, 2\}$ for $t = 1, 2, 3$ and $\varphi'(w_{a,1}') = \varphi'(w_{a,2}') \neq \varphi'(w_{a,3}')$.

**Proof:** Since the demand $\sigma'(D_{a,3}')$ is 5, $\varphi'(w_{a,3}')$ can be either of 1, 2, 3 or 4. However, it cannot be 3 or 4 since, in the same row as $w_{a,3}'$, there are 1-blocks in the subgrid $S_{a+1,1}$.
whose demands are 3 and 4. (See Figs. 11 and 12.) The latter can be confirmed easily.

Lemma 7: For any clause \( C_b \) \((b \in [M])\) with 2 literals, suppose that the positive literal is from the Boolean variable \( v_{a_1} \) and the negative literal is from \( v_{a_2} \). Let us denote by \( D'_1\) and \( D'_2\) the \((2 \times 1)\)-blocks in \( S_{a_1+1,1} \) and \( S_{a_2+1,1} \) that is in the same column as the cells \( w_{a_1,1} \) and \( w_{a_2,1} \) of \( U_b \). Then we have \( \varphi'(w_{a_1,1}) \neq \varphi'(w_{a_2,1}) \).

Proof: Since \( D'_1 \) and \( D'_2 \) are in the same column (see Figs. 7 and 11), \( \varphi'(w_{a_1,1}) \) and \( \varphi'(w_{a_2,1}) \) should not be equal from the Latin square condition.

Lemma 8: For any clause \( C_b \) \((b \in [M])\) with 3 literals, suppose that the positive literals are from the Boolean variables \( v_{a_1}, v_{a_2} \) and \( v_{a_3} \). Let us denote by \( D'_1\) and \( D'_2\) the \((2 \times 1)\)-block in \( S_{a_1+1,1} \) that is in the same column as the cell \( w_{a_1,1} \) (resp., \( w_{a_2,1} \) and \( w_{a_3,1} \) of \( U_b \). Then, \( \varphi'(w_{a_1,1}), \varphi'(w_{a_2,1}) \) and \( \varphi'(w_{a_3,1}) \) are not all equal.

Proof: See Figs. 8 and 11. The values \( \varphi'(w_{a_1,1}), \varphi'(w_{a_2,1}) \) and \( \varphi'(w_{a_3,1}) \) should not be all equal due to the three \((2 \times 1)\)-blocks that are in the 7th to 8th rows and in the 1st to 3rd columns in Fig. 8.

4. Discussion and Concluding Remarks

We showed that the decision problem version of BlockSum puzzle is NP-complete even if every block has size at most 2.

Although we focus on the existence of a solution of a given BlockSum instance, it is actually easy to generate an instance of BlockSum that has at least one solution; we can generate such an instance from a pair of an \( n \times n \) Latin square and a partition of the \( n \times n \) grid into blocks [8]. Thus, from the viewpoint of puzzle instance generation, it is important to consider the complexity of deciding whether a BlockSum instance has a solution other than the expected solution. This type of problem (and the complexity) is studied in terms of ASP-completeness, which does not require a polynomial time reduction from an NP-complete problem but requires an ASP-reduction from an ASP-complete problem [1]. Thus, it is an interesting and important open problem to decide whether BlockSum is ASP-complete or not.

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