A STRUCTURE THEOREM FOR STOCHASTIC PROCESSES
INDEXED BY THE DISCRETE HYPERCUBE

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ABSTRACT. Let $A$ be a finite set with $|A| \geq 2$, let $n$ be a positive integer, and let $A^n$ denote the discrete $n$-dimensional hypercube (that is, $A^n$ is the Cartesian product of $n$ many copies of $A$). Given a family $\langle D_t : t \in A^n \rangle$ of measurable events in a probability space (a stochastic process), what structural information can be obtained assuming that the events $\langle D_t : t \in A^n \rangle$ are not behaving as if they were independent? We obtain a complete answer to this problem (in a strong quantitative sense) subject to a mild “stationarity” condition. Our result has a number of combinatorial consequences, including a new (and the most informative so far) proof of the density Hales–Jewett theorem.

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1. INTRODUCTION

1.1. Motivation/Overview. Let $I$ be a nonempty finite set, let $\langle E_i : i \in I \rangle$ and $\langle D_i : i \in I \rangle$ be stochastic processes (families of measurable events) in a probability space $(\Omega, \mathcal{F}, \mathbb{P})$ with equal probability $\varepsilon > 0$, and assume that the events $\langle E_i : i \in I \rangle$ are independent. We wish to compare the distributions of the random variables

$$X = \sum_{i \in I} 1_{E_i} \quad \text{and} \quad Y = \sum_{i \in I} 1_{D_i}.$$
with the main goal being here that of transferring information from the distribution of $X$ (which we understand very well) to the distribution of $Y$, an object on which we have a priori no control. A classical method for doing so is by comparing the moments of $X$ and $Y$ (see, e.g., [Du]), a task which essentially reduces to that of comparing the joint probability of $\langle D_i : i \in F \rangle$ with the expected value $\varepsilon |F|$ as $F$ varies over all nonempty subsets of the index set $I$. Thus, assuming that the random variables $X$ and $Y$ are not close in distribution, then one is led to the following problem.

**Problem 1.1.** Let $F \subseteq I$ be nonempty, let $\sigma > 0$, and assume that

$$|P\left(\bigcap_{i \in F} D_i\right) - \varepsilon |F|| \geq \sigma.$$ 

What structural information can be obtained for the process $\langle D_i : i \in I \rangle$?

1.1.1. *The combinatorial content.* We will study Problem 1.1 in the case where the index set $I$ is a discrete hypercube, that is, a set of the form

$$A^n := A \times \cdots \times A$$

where $A$ is a finite set with $|A| \geq 2$ and $n$ is a positive integer which is commonly referred to as the dimension of the hypercube $A^n$. This choice of the index set is by no means arbitrary and it is ultimately related to the density Hales–Jewett theorem, a deep result due to Furstenberg and Katznelson [FK2] with numerous consequences in combinatorics, number theory, and theoretical computer science.

In order to properly discuss this relation we need to recall some basic definitions. Let $A$ and $n$ be as above, and fix a letter $x \notin A$ which we view as a variable. A variable word over $A$ of length $n$ is a finite sequence of length $n$ having values in $A \cup \{x\}$ where the letter $x$ appears at least once. If $v$ is a variable word over $A$ of length $n$ and $\alpha \in A$, then let $v(\alpha)$ denote the unique element of $A^n$ which is obtained by replacing every appearance of the letter $x$ in $v$ with $\alpha$. (For instance, if $A = \{\alpha, \beta, \gamma\}$ and $v = (\alpha, x, \gamma, \beta, x)$, then $v(\beta) = (\alpha, \beta, \gamma, \beta, \beta)$.) A combinatorial line of $A^n$ is a set of the form $\{v(\alpha) : \alpha \in A\}$ where $v$ is a variable word over $A$ of length $n$ (see [GRS, HJ]).

We are now in a position to recall the density Hales–Jewett theorem. We will state a probabilistic version—see, e.g., [FK2] Proposition 2.1—which is closer in spirit to our discussion. The relation between this probabilistic version and the more well-known combinatorial form which refers to dense subsets of discrete hypercubes will be discussed in Section 4.

**Theorem 1.2.** For every integer $k \geq 2$ and every $0 < \varepsilon \leq 1$ there exists a positive integer $\text{PHJ}(k, \varepsilon)$ with the following property. Let $A$ be a set with $|A| = k$, let $n \geq \text{PHJ}(k, \varepsilon)$ be an integer, and let $\langle D_t : t \in A^n \rangle$ be a stochastic process in a

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1This is because $X$ and $Y$ are both sums of indicator functions.
probability space $(\Omega, \mathcal{F}, P)$ such that $P(D_t) \geq \varepsilon$ for every $t \in A^n$. Then there exists a combinatorial line $L$ of $A^n$ such that

$$P\left(\bigcap_{t \in L} D_t\right) > 0.$$  

Of course, Theorem 1.2 is straightforward if the events $\langle D_t : t \in A^n \rangle$ are independent. Thus, the core of the theorem is to understand what happens when the events are not behaving as if they were independent, which is clearly an instance of Problem 1.1.

1.1.2. Deviating from the expected value: examples. To gain insight on the kind of structure one expects to obtain in Problem 1.1, it is useful to give examples of stochastic processes which exhibit non-independent behavior. Here and in the rest of this introduction, we will restrict our discussion on correlations over combinatorial lines. This is mainly because of the combinatorial importance of this case, but also because it is already quite representative of the behavior of correlations over more complicated sets.

Example 1.3. For concreteness we will work with the set $\{1, 2, 3\}$, but the argument can also be applied for any finite set $A$ with $|A| \geq 2$. Let $n$ be an arbitrary positive integer. We start with a family $\langle E_s : s \in \{1, 2\}^n \rangle$ of independent events in a probability space $(\Omega, \mathcal{F}, P)$ with equal probability $\varepsilon > 0$. Given $t \in \{1, 2, 3\}^n$ there are two natural ways to “project” it into $\{1, 2\}^n$. Specifically, let $t^{3 \rightarrow 1}$ and $t^{3 \rightarrow 2}$ denote the unique elements of $\{1, 2\}^n$ which are obtained by replacing every appearance of 3 in $t$ with 1 and 2 respectively. (E.g., if $t = (3, 2, 1, 3, 1) \in \{1, 2, 3\}^5$, then $t^{3 \rightarrow 1} = (1, 2, 1, 1, 1)$ and $t^{3 \rightarrow 2} = (2, 2, 1, 2, 1)$.) Then let $\langle D_t : t \in \{1, 2, 3\}^n \rangle$ be defined by setting $D_t := E^{t^{3 \rightarrow 1}} \cap E^{t^{3 \rightarrow 2}}$ for every $t \in \{1, 2, 3\}^n$.

Although the process $\langle D_t : t \in \{1, 2, 3\}^n \rangle$ in Example 1.3 is, arguably, quite easy to define, the analysis of its properties requires some work.

1.1.2.1. We first observe that for every $t \in \{1, 2, 3\}^n$ which contains 3 we have

$$P(D_t) = \varepsilon^2.$$  

Since the density of set of all elements of $\{1, 2, 3\}^n$ which do not contain 3 decreases exponentially with respect to the dimension $n$, we see that (1.2) holds true for “almost every” $t$.

1.1.2.2. The second basic property of the process $\langle D_t : t \in \{1, 2, 3\}^n \rangle$ concerns its correlations over combinatorial lines. Specifically, let $L = \{v(1), v(2), v(3)\}$ be a combinatorial line of $\{1, 2, 3\}^n$ where $v$ is a variable word over $\{1, 2, 3\}$ of length $n$ which contains 3. Then we have

$$P\left(\bigcap_{t \in L} D_t\right) = \varepsilon^4$$  

(1.3)
which implies that \( \langle D_t : t \in \{1, 2, 3\}^n \rangle \) exhibits non-independent\(^2\) behavior.

However, identity \((1.3)\) shows yet another important property of this process. More precisely, if \( v_1, v_2 \) are variable words over \( \{1, 2, 3\} \) of length \( n \) which both contain 3, then

\[
\Pr(D_{v_1(1)} \cap D_{v_1(2)} \cap D_{v_1(3)}) = \Pr(D_{v_2(1)} \cap D_{v_2(2)} \cap D_{v_2(3)}).
\]

(1.4)

In other words, the correlations of \( \langle D_t : t \in \{1, 2, 3\}^n \rangle \) over combinatorial lines are essentially constant. This property is abstracted in the following definition which originated\(^4\) in the work of Furstenberg and Katznelson [FK2].

**Definition 1.4 (Stationarity).** Let \( A \) be a finite set with \( |A| \geq 2 \), let \( n \) be a positive integer, let \( \eta > 0 \), and let \( \langle D_t : t \in A^n \rangle \) be a stochastic process in a probability space \((\Omega, \mathcal{F}, \Pr)\). We say that \( \langle D_t : t \in A^n \rangle \) is \( \eta \)-stationary (with respect to combinatorial lines) if for every nonempty \( \Gamma \subseteq A \) and every pair \( v_1, v_2 \) of variable words over \( A \) of length \( n \) we have

\[
\left| \Pr\left( \bigcap_{\alpha \in \Gamma} D_{v_1(\alpha)} \right) - \Pr\left( \bigcap_{\alpha \in \Gamma} D_{v_2(\alpha)} \right) \right| \leq \eta.
\]

(1.5)

(In particular, if \( \langle D_t : t \in A^n \rangle \) is an \( \eta \)-stationary process, then for every pair \( L_1, L_2 \) of combinatorial lines of \( A^n \) we have \( |\Pr(\bigcap_{t \in L_1} D_t) - \Pr(\bigcap_{t \in L_2} D_t)| \leq \eta \).)

Besides being very natural in this context\(^3\), stationarity is not a particularly restrictive condition. Indeed, it follows form a classical result due to Graham and Rothschild [GR] that stationary processes are the building blocks of arbitrary processes. (See Fact 3.1 in the main text.)

1.1.2.3. The last, and most significant, property of the process \( \langle D_t : t \in \{1, 2, 3\}^n \rangle \) is its hidden arithmetic structure which is described in the following definition.

**Definition 1.5 (Insensitivity).** Let \( A \) be a finite set with \( |A| \geq 2 \), let \( n \) be a positive integer, and let \( \alpha, \beta \in A \) with \( \alpha \neq \beta \).

1. Let \( s, t \in A^n \) and write \( s = (s_1, \ldots, s_n) \) and \( t = (t_1, \ldots, t_n) \). We say that \( s, t \) are \( (\alpha, \beta) \)-equivalent if for every \( i \in \{1, \ldots, n\} \) and every \( \gamma \in A \setminus \{\alpha, \beta\} \) we have that \( s_i = \gamma \) if and only if \( t_i = \gamma \). (Namely, \( s, t \) are \( (\alpha, \beta) \)-equivalent if they possibly differ only in the coordinates taking values in \( \{\alpha, \beta\} \).

2. We say that a stochastic process \( \langle D_t : t \in A^n \rangle \) in a probability space \((\Omega, \mathcal{F}, \Pr)\) is \( (\alpha, \beta) \)-insensitive provided that \( D_s = D_t \) for every \( s, t \in A^n \) which are \( (\alpha, \beta) \)-equivalent.

\(\)\(^2\)Specifically, by \((1.2)\), the expected probability in \((1.3)\) is \( \varepsilon^6 \).

\(\)\(^3\)The framework in [FK2] is somewhat different, but the essential content of Definition 1.4 is present in that work.

\(\)\(^4\)In particular note that, without assuming stationarity, one should instead study an averaged version of Problem 1.1.
The notion of insensitivity was introduced by Shelah in his proof of the Hales–Jewett theorem. It is the combinatorial analogue of the concept of a (discrete) Hilbert cube which is ubiquitous in additive combinatorics and arithmetic Ramsey theory (see, e.g., \cite{GRS, TV}).

Now, taking into account the definition of $t^3 \to 1$ and $t^3 \to 2$ in Example 1.3, it is easy to see that the processes \(E_{t^3 \to 1} : t \in \{1, 2, 3\}^n\) and \(E_{t^3 \to 2} : t \in \{1, 2, 3\}^n\) are \((1, 3)\)- and \((2, 3)\)-insensitive respectively. This property by itself yields that for every variable word \(v\) over \(\{1, 2, 3\}\) of length \(n\) we have

\[
D_{v(1)} \cap D_{v(2)} \cap D_{v(3)} = D_{v(1)} \cap D_{v(2)}.
\]

Note that identity (1.6) implies, in a rather extreme way, that the events \(D_{v(1)}, D_{v(2)}\) and \(D_{v(3)}\) cannot be independent. Thus we have a structural explanation of the fact that \(\langle D_t : t \in \{1, 2, 3\}^n \rangle\) exhibits non-independent behavior: it is the intersection of insensitive processes.

1.2. The main result. The following theorem (which is one of the main results of this paper and is proved in Section 3) shows that the example presented above is essentially the only example of a stationary process whose correlations over combinatorial lines deviate from what is expected.

**Theorem 1.6.** Let \(k \geq 2\) be an integer, and let \(\varepsilon, \sigma, \eta > 0\) such that

\[
\varepsilon \leq 1 - \frac{1}{2k}, \quad \sigma \leq \frac{\varepsilon^{k-1}}{2k} \quad \text{and} \quad \eta \leq \frac{\sigma}{4k-1}.
\]

Also let \(A\) be a set with \(|A| = k\), let \(n \geq k\) be an integer, and let \(\langle D_t : t \in A^n \rangle\) be an \(\eta\)-stationary process in a probability space \((\Omega, \mathcal{F}, \mathbb{P})\) such that \(|\mathbb{P}(D_t) - \varepsilon| \leq \eta\) for every \(t \in A^n\). Then, either

(i) for every combinatorial line \(L\) of \(A^n\) and every nonempty \(G \subseteq L\) we have

\[
|\mathbb{P}\left(\bigcap_{t \in G} D_t\right) - \varepsilon^{\left|G\right|}| \leq \sigma,
\]

(ii) or \(\langle D_t : t \in A^n \rangle\) correlates with a “structured” stochastic process; precisely, there exist a nonempty subset \(\Gamma\) of \(A\), \(\beta \in A \setminus \Gamma\) and a stochastic process \(\langle S_t : t \in A^n \rangle\) in \((\Omega, \mathcal{F}, \mathbb{P})\) such that the following are satisfied.

(a) For every \(t \in A^n\) we have \(S_t = \bigcap_{\alpha \in \Gamma} E_t^\alpha\) where for every \(\alpha \in \Gamma\) the process \(\langle E_t^\alpha : t \in A^n \rangle\) is \((\alpha, \beta)\)-insensitive.

(b) For every \(t \in A^n\) which contains \(\beta\) we have

\[
\mathbb{P}(S_t) \geq \frac{\varepsilon^{k-1}}{4k} \quad \text{and} \quad \mathbb{P}(D_t \mid S_t) \geq \varepsilon + \frac{\sigma}{4k-1}.
\]

\footnote{Actually, insensitivity was originally referring to subsets of discrete hypercubes (see Definition 1.5) and not to stochastic processes, but the difference between the two frameworks is minor.}

\footnote{This can be seen by identifying any nonempty finite set \(A\) with the interval \([1, \ldots, |A|]\) and then projecting the hypercube \(A^n\) into the integers via the map \(\langle \alpha_1, \ldots, \alpha_n \rangle \mapsto \sum_{i=1}^n \alpha_i |A|^{i-1}.\)
Theorem 1.6 is a new result whose most surprising feature is perhaps the fact that the conditional probability \( P(D_t \mid S_t) \) depends \textit{linearly} on the parameter \( \sigma \). As it is expected by Theorem 1.2, this information can in turn be used to prove the density Hales–Jewett theorem. We present this proof and we discuss in detail its quantitative aspects in Section 4. At this point we simply mention that it is a step towards obtaining primitive recursive bounds for the density Hales–Jewett numbers.

### 1.3. Correlations over arbitrary sets.

Beyond its combinatorial consequences, Theorem 1.6 is also the starting point of the analysis of correlations of stochastic processes over arbitrary nonempty subsets of discrete hypercubes. This analysis leads to a complete answer to Problem 1.1 and it is presented in the second part of this paper consisting of Sections 5–8. It can be seen as a natural—though not quite straightforward—generalization of the study of correlations over combinatorial lines. Specifically, there are two notable differences.

Firstly, the argument relies on the notion of the \textit{type}, a Ramsey-theoretic invariant which was introduced in \cite{DKT2} and encodes the “geometry” of a nonempty subset of a discrete hypercube. The definition of this invariant is recalled in Section 5, and it is crucially used in order to extend the notion of stationarity in this more general context (Definition 5.7 in the main text).

Secondly, the “structured” process which appears in part (ii.a) of Theorem 1.6 depends upon the type of the set \( G \) one is looking at part (i). This dependence is controlled by another invariant—the \textit{separation index}—which is introduced in Section 6. In particular, for correlations over sets which have the smallest possible separation index we have the exact analogue of Theorem 1.6 (Theorem 7.2 in the main text); however, the analogy breaks down at this point and the “structured” process which appears in part (ii.a) becomes more involved as the separation index increases (see Theorem 8.5 in the main text).

### 1.4. Outline of the argument.

The proof of Theorem 1.6 proceeds into two steps. In the first step and assuming that part (i) does not hold true, we select a subset \( B \) of \( A \) such that for every variable word \( v \) over \( A \) of length \( n \) and every nonempty proper subset \( \Sigma \) of \( B \) the events \( \langle D_{v(\alpha)} : \alpha \in \Sigma \rangle \) are essentially independent, yet the joint probability of \( \langle D_{v(\alpha)} : \alpha \in B \rangle \) deviates from the expected value. We emphasize that this selection is possible because the process \( \langle D_t : t \in A^n \rangle \) is stationary. The second step, which is the combinatorial heart of the matter, is to convert the irregularity of the correlations of \( \langle D_t : t \in A^n \rangle \) into correlation with a single structured process. This is achieved by taking advantage of the uniform

\footnote{For comparison, note that prior to this paper in order to obtain information as in part (ii.a) of Theorem 1.6 one needed to assume that for every combinatorial line \( L \) of \( A^n \) we have \( P(\bigcap_{t \in L} D_t) = 0 \); note that, because of Theorem 1.2 this assumption cannot hold true in the high-dimensional case.}

\footnote{The two parts are largely independent of each other and can be read separately.}
behavior of $\langle D_{\sigma(\alpha)} : \alpha \in B \rangle$ as $v$ varies over all variable words over $A$ of length $n$, and by carefully using the “projections” $i^3 \rightarrow 1$ and $i^3 \rightarrow 2$ described in Example 1.3 as well as their natural generalizations.

The argument for the case of correlations over arbitrary sets follows the same outline, though the details are—as expected—more complicated. We comment on the differences of the proof of the general case in Sections 7 and 8.

2. Combinatorial background

2.1. By $\mathbb{N} = \{0, 1, 2, \ldots \}$ we denote the set of all natural numbers, and for every positive integer $n$ we set $[n] := \{1, \ldots, n\}$. For every set $X$ by $|X|$ we denote its cardinality; moreover, for every subset $A$ of $X$ by $A^c$ we denote the complement of $A$, that is, $A^c := X \setminus A$.

2.2. Definitions. Let $A$ denote a finite set with $|A| \geq 2$.

2.2.1. As in [1], for every positive integer $n$ by $A^n$ we denote the Cartesian product of $n$ many copies of $A$; we view $A^n$ as the set of all sequences of length $n$ having values in $A$. Also let $\emptyset$ denote the empty sequence, set $A^0 := \{\emptyset\}$, and let

$$ A^{\leq n} := \bigcup_{n \in \mathbb{N}} A^n $$

denote the set of all finite (possibly empty) sequences in $A$. For every $t, s \in A^{\leq n}$ by $t \cdot s$ we denote the concatenation of $t$ and $s$; notice, in particular, that if $t \in A^n$ and $s \in A^m$ for some $n, m \in \mathbb{N}$, then $t \cdot s \in A^{n+m}$.

2.2.2. Variable words. Let $n, m$ be positive integers, and fix a set $\{x_1, \ldots, x_m\}$ which is disjoint from $A$; we view $\{x_1, \ldots, x_m\}$ as a set of variables. An $m$-variable word over $A$ of length $n$ is a finite sequence $v$ of length $n$ having values in $A \cup \{x_1, \ldots, x_m\}$ such that: (1) for every $i \in [m]$ the letter $x_i$ appears in $v$ at least once, and (2) if $m \geq 2$, then for every $i, j \in [m]$ with $i < j$ all appearances of $x_i$ in $v$ precede all appearances of $x_j$. If $v$ is an $m$-variable word over $A$ of length $n$ and $\alpha_1, \ldots, \alpha_m \in A$, then by $v(\alpha_1, \ldots, \alpha_m)$ we denote the unique element of $A^n$ which is obtained by replacing every appearance of $x_i$ in $v$ with $\alpha_i$ for every $i \in [m]$. (For example, if $A = \{\alpha, \beta, \gamma\}$ and $v = (x_1, \gamma, x_2, x_2, \beta, x_3)$, then $v(\beta, \alpha, \gamma) = (\beta, \gamma, \alpha, \alpha, \beta, \gamma)$.)

2.2.3. Combinatorial spaces and canonical isomorphisms. A combinatorial space of $A^{\leq n}$ is a subset $V$ of $A^{\leq n}$ of the form

$$ V = \{v(\alpha_1, \ldots, \alpha_m) : \alpha_1, \ldots, \alpha_m \in A\} $$

where $m$ is a positive integer and $v$ is an $m$-variable word over $A$ of length $n$ for some positive integer $n$ (in particular, we have $V \subseteq A^n$). Notice that $m, v$ and $n$ are unique since $|A| \geq 2$; the (unique) positive integer $m$ is called the dimension of $V$ and is denoted by $\dim(V)$. Also observe that the 1-dimensional combinatorial spaces are precisely the combinatorial lines already mentioned in the introduction.
Finally, if $V_1$ and $V_2$ are two combinatorial spaces of $A^{<N}$, then we say that $V_1$ is a **combinatorial subspace** of $V_2$ provided that $V_1 \subseteq V_2$.

We view an $m$-dimensional combinatorial space $V$ as a “copy” of $A^m$ inside $A^{<N}$, and we will identify $V$ with $A^m$ for most practical purposes. To this end, we introduce the following definition.

**Definition 2.1.** Let $A$ be a finite set with $|A| \geq 2$, and let $V$ be a combinatorial space of $A^{<N}$. Set $m := \dim(V)$ and let $v$ be the unique $m$-variable word over $A$ which generates $V$ via formula (2.2). The canonical isomorphism associated with $V$ is the bijection $I_V : A^m \to V$ defined by the rule
\[
I_V((\alpha_1, \ldots, \alpha_m)) = v(\alpha_1, \ldots, \alpha_m).
\]
for every $(\alpha_1, \ldots, \alpha_m) \in A^m$.

Note that canonical isomorphisms preserve combinatorial subspaces and their dimension; precisely, if $V$ is an $m$-dimensional combinatorial space of $A^{<N}$ and $W \subseteq A^m$, then $W$ is a combinatorial subspace of $A^m$ with $\dim(W) = \ell$ if and only if $I_V(W)$ is a combinatorial subspace of $V$ with $\dim(I_V(W)) = \ell$. For an exposition of the properties of canonical isomorphisms we refer to [DK, Section 1.3].

**2.3. Colorings of combinatorial lines.** We will need the following special case of the Graham–Rothschild theorem [GR]. The corresponding primitive recursive bounds are taken from [Ty].

**Proposition 2.2.** For every triple $k, m, r$ of positive integers with $k \geq 2$ there exists a positive integer $N$ with the following property. For every set $A$ with $|A| = k$, every combinatorial space $V$ of $A^{<N}$ with $\dim(V) \geq n$ and every $r$-coloring of the set of all combinatorial lines of $V$ there exists an $m$-dimensional combinatorial subspace $W$ of $V$ such that the set of all combinatorial lines of $W$ is monochromatic. The least positive integer $N$ with this property is denoted by $GRL(k, m, r)$.

Moreover, the numbers $GRL(k, m, r)$ are upper bounded by a primitive recursive function belonging to the class $E^5$ of Grzegorczyk’s hierarchy.

For a discussion of Grzegorczyk’s hierarchy of primitive recursive functions and its role in analyzing the bounds associated with various results in Ramsey theory we refer to [DK] Appendix A.

3. Correlations over combinatorial lines

In this section we give the proof of Theorem 1.6. As we have noted in the introduction, the argument (which also pertains the proofs of Theorems 7.2 and 8.5) can be roughly summarized by saying that higher order correlations of a process

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9Actually, the Graham–Rothschild theorem refers to parameter words, a concept which is slightly different from the notion of a variable word. However, for colorings of combinatorial lines the difference between the two frameworks is minor.
can be converted into correlation with a single structured process. Perhaps the most transparent instance of this fact is the proof of Proposition 3.7 below.

We begin with some preliminary steps, including a discussion on some basic properties of stationary processes.

3.1. Stationarity. We have already noted that the Graham–Rothschild theorem (more precisely, Proposition 2.2) implies that stationary processes are the building blocks of arbitrary processes. In particular, we have the following fact. The proof is straightforward.

**Fact 3.1.** Let \( k \geq 2 \) be an integer, and let \( A \) be a set with \( |A| = k \). Also let \( 0 < \eta \leq 1 \), and let \( n, m \) be positive integers such that
\[
\eta \leq GRL(k, m, [1/\eta]^{2^k-1}).
\]
Then for every stochastic process \( \langle D_t : t \in A^n \rangle \) in a probability space \((\Omega, \mathcal{F}, \mathbb{P})\) there exists an \( m \)-dimensional combinatorial subspace \( V \) of \( A^n \) such that the process \( \langle D_{t\Gamma(s)} : s \in A^m \rangle \) (namely, the restriction of \( \langle D_t : t \in A^n \rangle \) on \( V \)) is \( \eta \)-stationary.

The following lemma shows that one can upgrade the estimate in (1.5) and stabilize the joint distribution of certain boolean combinations of the events of a stationary processes. (Here, and in the rest of this paper, we follow that convention that the intersection of an empty family of events of a probability space \((\Omega, \mathcal{F}, \mathbb{P})\) is equal to the sample space \( \Omega \).)

**Lemma 3.2.** Let \( A \) be a finite set with \( |A| \geq 2 \), let \( n \) be a positive integer, let \( \eta > 0 \), and let \( \langle D_t : t \in A^n \rangle \) be an \( \eta \)-stationary stochastic process in a probability space \((\Omega, \mathcal{F}, \mathbb{P})\). Then for every pair \( \Gamma_1, \Gamma_2 \subseteq A \) with \( \Gamma_1 \cap \Gamma_2 = \emptyset \) and every pair \( v_1, v_2 \) of variable words over \( A \) of length \( n \) we have
\[
\left| \mathbb{P}\left( \bigcap_{a \in \Gamma_1} D_{v_1(a)} \cap \bigcap_{a \in \Gamma_2} D_{v_2(a)}^c \right) - \mathbb{P}\left( \bigcap_{a \in \Gamma_1} D_{v_2(a)} \cap \bigcap_{a \in \Gamma_2} D_{v_2(a)}^c \right) \right| \leq 2^{\|\Gamma_1\|} \eta.
\]

**Proof.** Let \( \Gamma_1, \Gamma_2, v_1, v_2 \) be as in the statement of the lemma. Then, using the inclusion–exclusion formula, we have
\[
\left| \mathbb{P}\left( \bigcap_{a \in \Gamma_1} D_{v_1(a)} \cap \bigcap_{a \in \Gamma_2} D_{v_2(a)}^c \right) - \mathbb{P}\left( \bigcap_{a \in \Gamma_1} D_{v_2(a)} \cap \bigcap_{a \in \Gamma_2} D_{v_2(a)}^c \right) \right|
= \left| \mathbb{P}\left( \bigcap_{a \in \Gamma_1} D_{v_1(a)} \cap \bigcup_{a \in \Gamma_2} D_{v_2(a)}^c \right) - \mathbb{P}\left( \bigcap_{a \in \Gamma_1} D_{v_2(a)} \cap \bigcup_{a \in \Gamma_2} D_{v_2(a)}^c \right) \right|
\leq \sum_{\Gamma_2' \subseteq \Gamma_2} \left| \mathbb{P}\left( \bigcap_{a \in \Gamma_1} D_{v_1(a)} \cap \bigcap_{a \in \Gamma_2'} D_{v_1(a)} \right) - \mathbb{P}\left( \bigcap_{a \in \Gamma_1} D_{v_2(a)} \cap \bigcap_{a \in \Gamma_2'} D_{v_2(a)} \right) \right| \leq 2^{\|\Gamma_2\|} \eta
\]
and the proof is completed. \(\square\)

**Remark 3.3.** We notice that the assumption in Theorem 1.6 that \( |\mathbb{P}(D_t) - \varepsilon| \leq \eta \) for every \( t \in A^n \) follows from \( \eta \)-stationarity provided that the dimension \( n \) is sufficiently large. Indeed, let \( A, n \) and \( \langle D_t : t \in A^n \rangle \) be as in Theorem 1.6 clearly,
we have $n \geq |A|$. We select $t_0 \in A^n$ such that every $\alpha \in A$ appears in $t_0$ at least once (this selection is possible since $n \geq |A|$), and we set $\varepsilon := \max\{\mathbb{P}(D_{t_0}), \eta\} > 0$. Note that for every $t \in A^n$ there exist two variable words $v_1, v_2$ over $A$ of length $n$ and $\alpha \in A$ such that $t = v_1(\alpha)$ and $t_0 = v_2(\alpha)$. Invoking (1.5), we conclude that $|\mathbb{P}(D_t) - \varepsilon| \leq \eta$.

3.2. Insensitivity. Let $A$ be a finite set with $|A| \geq 2$, let $n$ be a positive integer, and let $\alpha, \beta \in A$ with $\alpha \neq \beta$. As in Example 1.3, for every $t \in A^n$ let $t^{\beta \rightarrow \alpha}$ denote the unique element of $(A \setminus \{\beta\})^n$ which is obtained by replacing every appearance of $\beta$ in $t$ with $\alpha$. We will use this operation in order to produce insensitive processes. To this end, we will need the following elementary (though crucial) fact. Its proof is straightforward.

**Fact 3.4.** Let $A$ be a finite set with $|A| \geq 2$, let $n$ be a positive integer, and let $\alpha, \beta \in A$ with $\alpha \neq \beta$. Then the map $A^n \ni t \mapsto t^{\beta \rightarrow \alpha} \in (A \setminus \{\beta\})^n$ is a projection; that is, for every $t \in A^n$ which does not contain $\beta$ we have that $t^{\beta \rightarrow \alpha} = t$. Moreover, if $t, s \in A^n$ are $(\alpha, \beta)$-equivalent, then $t^{\beta \rightarrow \alpha} = s^{\beta \rightarrow \alpha}$.

3.3. Pseudorandomness, supercorrelation, subcorrelation. Let $E_1, \ldots, E_\ell$ be measurable events in a probability space with equal probability $\varepsilon > 0$. Notice that the joint probability of $E_1, \ldots, E_\ell$ can be naturally categorized according to whether it is greater than, less than, or almost equal to the expected value $\varepsilon^\ell$. As expected, our analysis depends on this trichotomy, and as such, it is convenient to introduce the following definition.

**Definition 3.5.** Let $A$ be a finite set with $|A| \geq 2$, let $n \geq |A|$ be an integer, let $0 < \eta, \varepsilon \leq 1$, and let $\langle D_t : t \in A^n \rangle$ be an $\eta$-stationary process in a probability space $(\Omega, \mathcal{F}, \mathbb{P})$ such that $|\mathbb{P}(D_t) - \varepsilon| \leq \eta$ for every $t \in A^n$. Also let $\Gamma \subseteq A$ be nonempty, and let $\theta \geq 0$.

1. (Pseudorandomness) We say that $\langle D_t : t \in A^n \rangle$ is $(\Gamma, \theta)$-pseudorandom if $|\mathbb{P}(\bigcap_{t \in F} D_v(\alpha)) - \varepsilon|^{1/|F|} \leq \theta$ for every variable word $v$ over $A$ of length $n$.
2. (Supercorrelation) We say that $\langle D_t : t \in A^n \rangle$ is $(\Gamma, \theta)$-supercorrelated if $\mathbb{P}(\bigcap_{t \in F} D_v(\alpha)) \geq \varepsilon^{1/|F|} + \theta$ for every variable word $v$ over $A$ of length $n$.
3. (Subcorrelation) We say that $\langle D_t : t \in A^n \rangle$ is $(\Gamma, \theta)$-subcorrelated if $\mathbb{P}(\bigcap_{t \in F} D_v(\alpha)) \leq \varepsilon^{1/|F|} - \theta$ for every variable word $v$ over $A$ of length $n$.

We have the following fact.

**Fact 3.6.** Let $A, n, \eta, \varepsilon$ and $\langle D_t : t \in A^n \rangle$ be as in Definition 3.5. Also let $\Gamma \subseteq A$ be nonempty, and let $\theta \geq \eta$. Then one of the following holds true.

1. The process $\langle D_t : t \in A^n \rangle$ is $(\Gamma, \theta)$-pseudorandom.
2. The process $\langle D_t : t \in A^n \rangle$ is $(\Gamma, \theta - \eta)$-supercorrelated.
3. The process $\langle D_t : t \in A^n \rangle$ is $(\Gamma, \theta - \eta)$-subcorrelated.
Proof. Assume that (i) does not hold true, that is, there is a variable word \( v \) over \( A \) of length \( n \) such that either \( \mathbb{P}(\bigcap_{\alpha \in \Gamma} D_{\nu(\alpha)}) \geq \varepsilon|\Gamma| + \theta \), or \( \mathbb{P}(\bigcap_{\alpha \in \Gamma} D_{\nu(\alpha)}) \leq \varepsilon|\Gamma| - \theta \). Invoking the \( \eta \)-stationarity of \( \langle D_t : t \in A^n \rangle \), we see that the first alternative yields part (ii), while the second alternative yields part (iii).

We are ready to state the main result in this subsection.

**Proposition 3.7.** Let \( A \) be a finite set with \( |A| \geq 2 \), and let \( n \geq |A| \) be an integer. Also let \( 0 < \eta, \varepsilon \leq 1 \), and let \( \langle D_t : t \in A^n \rangle \) be a \( \eta \)-stationary stochastic process in a probability space \( (\Omega, \mathcal{F}, \mathbb{P}) \) such that \( |\mathbb{P}(D_t) - \varepsilon| \leq \eta \) for every \( t \in A^n \). Finally, let \( \theta, \sigma \geq 0 \), let \( \Gamma \subseteq A \) be nonempty, and let \( \beta \in A \setminus \Gamma \). Assume that \( \langle D_t : t \in A^n \rangle \) is \((\Gamma, \theta)\)-pseudorandom, and set \( p := |\Gamma| \). Then there exists a stochastic process \( \langle S_t : t \in A^n \rangle \) in \( (\Omega, \mathcal{F}, \mathbb{P}) \) with the following properties.

\( \begin{align*} 
\text{(i)} & \quad \text{For every } t \in A^n \text{ we have } S_t = \bigcap_{\alpha \in \Gamma} E^\alpha_t \text{ where for every } \alpha \in \Gamma \text{ the process } \langle E^\alpha_t : t \in A^n \rangle \text{ is } (\alpha, \beta)\text{-insensitive.} \\
\text{(ii)} & \quad \text{For every } t \in A^n \text{ which does not contain } \beta \text{ and every } \alpha \in \Gamma \text{ we have } E^\alpha_t = D_t. \text{ (Thus, } S_t = D_t \text{ for every } t \in A^n \text{ which does not contain } \beta.) \\
\text{(iii)} & \quad \text{For every } t \in A^n \text{ which contains } \beta \text{ we have } |\mathbb{P}(S_t) - \mathbb{P}(\beta)| \leq \theta. \\
\text{(iv)} & \quad \text{If } \langle D_t : t \in A^n \rangle \text{ is } (\Gamma \cup \{\beta\}, \sigma)\text{-supercorrelated, then for every } t \in A^n \text{ which contains } \beta \text{ we have} \\
& \quad \quad |\mathbb{P}(D_t | S_t) - \varepsilon| \leq \varepsilon \left(1 + \frac{\sigma \varepsilon^{-1} - \theta}{\varepsilon \sigma - \theta} \right). \\
\text{(v)} & \quad \text{If } \langle D_t : t \in A^n \rangle \text{ is } (\Gamma \cup \{\beta\}, \sigma)\text{-subcorrelated, then for every } t \in A^n \text{ which contains } \beta \text{ we have} \\
& \quad \quad |\mathbb{P}(D_t | S_t) - \varepsilon| \leq \varepsilon \left(1 - \frac{\sigma \varepsilon^{-1} - \theta}{\varepsilon \sigma - \theta} \right). \\
\end{align*} \)

Proof. We first observe that the conditions in parts (i) and (ii) completely determine the stochastic process \( \langle E^\alpha_t : t \in A^n \rangle \) for every \( \alpha \in \Gamma \). However, it is possible to give an alternative (and more intrinsic) definition of these processes which facilitates the proofs of parts (iii)–(v) and it is easier to generalize when considering correlations over more complicated sets (see, in particular, Sections 7 and 8). More precisely, notice that, by Fact 3.4 we have \( E^\alpha_t = D_{\beta \rightarrow \alpha} \) for every \( t \in A^n \) and every \( \alpha \in \Gamma \). We will also need the following important property of this construction. For every \( t \in A^n \) which contains \( \beta \) let \( v_t \) denote the unique variable word over \( A \setminus \{\beta\} \) of length \( n \) which is obtained by replacing every appearance of \( \beta \) in \( t \) with the variable \( x \), and note that \( t = v_t(\beta) \) and \( t^\beta \rightarrow ^\alpha = v_t(\alpha) \) for every \( \alpha \in \Gamma \). Consequently, for every \( t \in A^n \) which contains \( \beta \) we have

\( \begin{align*} 
\text{(3.5)} & \quad S_t = \bigcap_{\alpha \in \Gamma} D_{v_t(\alpha)} \quad \text{and} \quad D_t \cap S_t = \bigcap_{\alpha \in \Gamma \cup \{\beta\}} D_{v_t(\alpha)}. \\
\end{align*} \)

After this preliminary discussion, we are ready to proceed to the rest of the proof. Part (iii) follows immediately by the first identity in (3.5) and our assumption that the process \( \langle D_t : t \in A^n \rangle \) is \((\Gamma, \theta)\)-pseudorandom.
For part (iv), assume that $\langle D_t : t \in A^n \rangle$ is $(\Gamma \cup \{\beta\}, \sigma)$-super correlated. Fix $t \in A^n$ which contains $\beta$. By the second identity in (3.5) and the supercorrelation assumption, we see that $\mathbb{P}(D_t \cap S_t) \geq \varepsilon^{p+1} + \sigma$; on the other hand, by part (iii), we have $\mathbb{P}(S_t) \leq \varepsilon^p + \theta$. Therefore,

$$\mathbb{P}(D_t \mid S_t) = \frac{\mathbb{P}(D_t \cap S_t)}{\mathbb{P}(S_t)} \geq \frac{\varepsilon^{p+1} + \sigma}{\varepsilon^p + \theta} = \varepsilon \left(1 + \frac{\sigma \varepsilon^{-1} - \theta}{\varepsilon^p + \theta}\right)$$

as desired.

Finally, assume that $\langle D_t : t \in A^n \rangle$ is $(\Gamma \cup \{\beta\}, \sigma)$-sub correlated, and fix $t \in A^n$ which contains $\beta$. As above, using the second identity in (3.5) and the subcorrelation assumption, we obtain that $\mathbb{P}(D_t \cap S_t) \leq \varepsilon^{p+1} - \sigma$. By part (iii), we have $\mathbb{P}(S_t) \geq \varepsilon^p - \theta$, and so,

$$\mathbb{P}(D_t \mid S_t) = \frac{\mathbb{P}(D_t \cap S_t)}{\mathbb{P}(S_t)} \leq \frac{\varepsilon^{p+1} - \sigma}{\varepsilon^p - \theta} = \varepsilon \left(1 - \frac{\sigma \varepsilon^{-1} - \theta}{\varepsilon^p - \theta}\right).$$

The proof is completed. \qed

Remark 3.8. Observe that the variable word $v_t$ defined in the proof of Proposition 3.7 is not typical since it does not contain $\beta$. Nevertheless, because stationarity is a global property, it is possible to have information for the correlation of the events $\langle D_{v_t(\alpha)} : \alpha \in \Gamma \rangle$. This fact (namely, the necessity to understand the correlations of $\langle D_t : t \in A^n \rangle$ over sparse sets of combinatorial lines) is rather subtle and appears to be a genuine obstacle for extending Theorem 1.6 to not necessarily stationary processes.

Remark 3.9 (Extreme cases). Note that the extreme cases in Proposition 3.7 are: (a) “$\theta = 0$” and “$\sigma = \varepsilon^p - \varepsilon^{p+1}$” if the stochastic process $\langle D_t : t \in A^n \rangle$ is super correlated, and (b) “$\theta = 0$” and “$\sigma = \varepsilon^{p+1}$” if $\langle D_t : t \in A^n \rangle$ is sub correlated. In the first case we have that $\mathbb{P}(D_t \mid S_t) = 1$ for every $t \in A^n$ containing $\beta$, which is clearly equivalent to saying that $S_t \subseteq D_t$. Examples of stochastic processes of this form can be obtained by modifying (in a straightforward way) Example 1.3. At the other extreme, we see that $\mathbb{P}(D_t \mid S_t) = 0$ for every $t \in A^n$ which contains $\beta$. In contrast to the previous case, this phenomenon cannot occur if the dimension $n$ is sufficiently large; this is a consequence of Theorem 1.2.

3.4. **Proof of Theorem 1.6**. We begin by introducing a finite sequence $(\theta_p)_{p=0}^k$ of positive reals defined by the rule

$$\left\{
\begin{array}{l}
\theta_0 = 0, \, \theta_1 = \eta, \\
\theta_p = 4p^{-k} \sigma \text{ if } p \in \{2, \ldots, k\}.
\end{array}\right. \tag{3.6}$$

(Note that, by (1.7), the sequence $(\theta_p)_{p=0}^k$ is increasing.) Next observe that if for every nonempty $\Gamma \subseteq A$ the process $\langle D_t : t \in A^n \rangle$ is $(\Gamma, \theta_{|\Gamma})$-pseudorandom, then part (i) of the theorem holds true. Therefore, we may assume that there exists nonempty $\Delta \subseteq A$ such that $\langle D_t : t \in A^n \rangle$ is not $(\Delta, \theta_{|\Delta})$-pseudorandom. We fix a nonempty subset $\Gamma_0$ of $A$ which satisfies this property and with minimal
cardinality. (Notice, in particular, that if \( \Sigma \) is a nonempty proper subset of \( \Gamma_0 \),
then \( \langle D_t : t \in A^n \rangle \) is \((\Sigma, \theta_{\Sigma})\)-pseudo-random.) Since the process \( \langle D_t : t \in A^n \rangle \)
is \( \eta \)-stationary and \( \theta_1 = \eta \), we see that \(|\Gamma_0| \geq 2 \). We select \( \beta \in \Gamma_0 \), and we set
\( \Gamma := \Gamma_0 \setminus \{ \beta \} \) and \( p := |\Gamma| \); observe that \( 1 \leq p \leq k - 1 \). Set \( \theta := \theta_p \) and \( \Theta := \theta_{p+1} \).

By Fact 3.6 and our assumption that the stochastic process \( \langle D_t : t \in A^n \rangle \) is not
\((\Gamma \cup \{ \beta \}, \Theta)\)-pseudo-random, we see that either

\( \text{Case 1:} \) \( \langle D_t : t \in A^n \rangle \) is \((\Gamma \cup \{ \beta \}, \Theta - \eta)\)-subcorrelated,

\( \text{Case 2:} \) \( \langle D_t : t \in A^n \rangle \) is \((\Gamma \cup \{ \beta \}, \Theta - \eta)\)-subcorrelated.

We will show that in both cases part (ii) of the theorem holds true.

\( \text{Case 1:} \) \( \langle D_t : t \in A^n \rangle \) is \((\Gamma \cup \{ \beta \}, \Theta - \eta)\)-subcorrelated. By Proposition 3.7
applied for “\( \sigma = \Theta - \eta \)”, there exists a process \( \langle S_t : t \in A^n \rangle \) which satisfies part (ii.a) of the theorem such that for every \( t \in A^n \) which contains \( \beta \) we have

\[
\begin{align*}
\text{(a)} & \quad |P(S_t) - \varepsilon^n| \leq \theta, \\
\text{(b)} & \quad P(D_t | S_t) \geq \varepsilon (1 + \frac{\Theta - \eta - \varepsilon \theta}{\varepsilon + \theta} - \frac{\varepsilon k^{-1}}{4k}).
\end{align*}
\]

Therefore, by (a) above and the fact that \( \theta \leq \sigma/4 \), for every \( t \in A^n \) which contains \( \beta \) we have

\[
P(S_t) \geq \varepsilon^n - \theta \geq \varepsilon^{k-1} - \frac{\varepsilon k^{-1}}{4k} \geq \frac{\varepsilon k^{-1}}{4k}
\]

while, by (b) and the fact that \( \eta, \theta \leq \Theta/4 \),

\[
P(D_t | S_t) \geq \varepsilon + \Theta - \eta - \varepsilon \theta \geq \varepsilon + \frac{1}{2}(\Theta - \eta - \theta) \geq \varepsilon + \frac{\Theta}{4} \geq \varepsilon + \frac{\Theta}{4} = \varepsilon + \frac{\sigma}{4k - 1}.
\]

Thus, part (ii.b) of the theorem is also satisfied, as desired.

\( \text{Case 2:} \) \( \langle D_t : t \in A^n \rangle \) is \((\Gamma \cup \{ \beta \}, \Theta - \eta)\)-subcorrelated. For every \( t \in A^n \) and
every \( \alpha \in \Gamma \) set \( E_t^\alpha := D_t^{\beta \rightarrow \alpha} \). (Recall that \( t^{\beta \rightarrow \alpha} \) denotes the unique element of
\((A \setminus \{ \beta \})^n \) which is obtained by replacing every appearance of \( \beta \) in \( t \) with \( \alpha \).) We
select \( \gamma \in \Gamma \), we set \( B := \Gamma \setminus \{ \gamma \} \), and for every \( t \in A^n \) we define

\[
S_t := \left( \bigcap_{\alpha \in B} E_t^\alpha \right) \cap (E_t^\gamma)^C.
\]

(Recall that, by convention, \( \bigcap_{\alpha \in B} E_t^\alpha = \Omega \) if \( B = \emptyset \).) Clearly, the stochastic process \( \langle S_t : t \in A^n \rangle \) satisfies part (ii.a) of the theorem. As in the proof of Proposition 3.7
for every \( t \in A^n \) which contains \( \beta \) by \( v_t \) we denote the unique variable word over
\( A \setminus \{ \beta \} \) of length \( n \) which is obtained by replacing every appearance of \( \beta \) in \( t \) with the
variable \( x \); recall that \( t = v_t(\beta) \) and \( t^{\beta \rightarrow \alpha} = v_t(\alpha) \) for every \( \alpha \in \Gamma \). Consequently, for every \( t \in A^n \) which contains \( \beta \) we have

\[
S_t = \left( \bigcap_{\alpha \in B} D_{v_t(\alpha)} \right) \setminus \left( \bigcap_{\alpha \in \Gamma} D_{v_t(\alpha)} \right)
\]
and
\[(3.9) \quad D_t \cap S_t = \left( \bigcap_{\alpha \in B \cup \{\beta\}} D_{v_1(\alpha)} \right) \setminus \left( \bigcap_{\alpha \in \Gamma \cup \{\beta\}} D_{v_1(\alpha)} \right).\]

Since \(\langle D_t : t \in A^n \rangle\) is
- \((\Gamma \cup \{\beta\}, \Theta - \eta)\)-subcorrelated,
- \((B, \theta_{p-1})\)-pseudorandom if \(p > 1\) (if \(p = 1\), then this is superfluous), and
- \((\Gamma, \theta)\)-pseudorandom and \((B \cup \{\beta\}, \theta)\)-pseudorandom,

for every \(t \in A^n\) which contains \(\beta\) we have
\[
\mathbb{P}(S_t) \leq (\varepsilon^{p-1} + \theta_{p-1}) - (\varepsilon^p - \Theta) \quad \text{and} \quad \mathbb{P}(D_t \cap S_t) \geq (\varepsilon^p - \Theta) - (\varepsilon^{p+1} - \Theta + \eta).
\]

Moreover, by \((1.7)\) and \((3.6)\), we have \(\theta + \theta_{p-1} \leq \varepsilon^p, \theta \leq \Theta/4\) and \(\theta_{p-1} + \eta \leq \Theta/4\).

Therefore, for every \(t \in A^n\) which contains \(\beta\)
\[
\mathbb{P}(D_t | S_t) \geq \varepsilon^p - \varepsilon^{p+1} + \Theta - \theta - \eta \geq \varepsilon + \Theta - 2\theta - \theta_{p-1} - \eta \geq \varepsilon^p - \varepsilon^p + \theta + \theta_{p-1} \geq \varepsilon + \Theta + \frac{\sigma}{4^{k-1}}.
\]

Finally, by \((1.7)\), \((3.8)\) and the fact that \(\langle D_t : t \in A^n \rangle\) is \((B, \theta_{p-1})\)-pseudorandom and \((\Gamma, \theta)\)-pseudorandom, we conclude that
\[
\mathbb{P}(S_t) \geq \varepsilon^{p-1} - \varepsilon^p - \theta - \theta_{p-1} \geq \varepsilon^{k-1}(1 - \varepsilon) - 2\theta \geq \frac{\varepsilon^{k-1}}{2k} - 2\theta \geq \frac{\varepsilon^{k-1}}{4k}
\]

for every \(t \in A^n\) which contains \(\beta\). The proof is completed.

4. PROOF OF THE DENSITY HALEES–JEWETT THEOREM

4.1. In this section we give a proof of the density Hales–Jewett theorem which is based on Theorem 1.6. We begin by recalling the combinatorial version of the density Hales–Jewett theorem. (The reader is advised to compare this version with Theorem 1.2 stated in the introduction.)

**Theorem 4.1.** For every integer \(k \geq 2\) and every \(0 < \delta \leq 1\) there exists a positive integer \(\text{DHJ}(k, \delta)\) with the following property. Let \(A\) be a set with \(|A| = k\), and let \(n \geq \text{DHJ}(k, \delta)\) be an integer. Then every \(D \subseteq A^n\) with \(|D| \geq \delta |A^n|\) contains a combinatorial line of \(A^n\).

There are several effective proofs\(^{10}\) of Theorem 4.1; see [DKT1, P2, Tao]. Despite this progress, the understanding of the behavior of the density Hales–Jewett numbers \(\text{DHJ}(k, \delta)\) is rather poor. Indeed, the best known upper bounds are obtained in [P2] and have an Ackermann-type dependence with respect to \(k\).

The proof of Theorem 4.1 given below is based on a density increment strategy (a method introduced by Roth [Roth]) and follows the general scheme developed in [P2]. Its most important feature is the quantitative improvement of a crucial

\(^{10}\) Another ergodic-theoretic proof was given in [Am].
step which appears (in various forms) in all known combinatorial proofs of the density Hales–Jewett theorem. (We discuss this particular feature in Remark 4.7 below.) The driving force behind this improvement is Theorem 1.6.

4.1.1. Step 1: from dense subsets of discrete hypercubes to stochastic processes. Strictly speaking, this step is not an internal part of the proof of Theorem 4.1. However, it is conceptually significant since it enables us to pass from dense sets to stochastic processes. This is essentially the content of the following simple lemma whose proof can be found in [DKT1, Lemma 4].

Lemma 4.2. Let \( k, m \) be positive integers with \( k \geq 2 \), let \( 0 < \eta \leq 1 \), let \( A \) be a set with \( |A| = k \), and let \( n \) be a positive integer such that

\[
(4.1) \quad n \geq \frac{km}{\eta}.
\]

Then for every \( D \subseteq A^n \) there exist \( t \in \{m, \ldots, n-1\} \) and an \( m \)-dimensional combinatorial subspace \( V \) of \( A^t \) such that for every \( t \in V \) we have

\[
(4.2) \quad \frac{|D_t|}{|A^{n-t}|} \geq \frac{|D|}{|A^n|} - \eta
\]

where \( D_t = \{ s \in A^{n-t} : t^-s \in D \} \) denotes the section of \( D \) at \( t \).

Remark 4.3. There is a more powerful probabilistic version of Lemma 4.2 which can be stated as a concentration inequality and relies on properties of martingale difference sequences; see [DKT3, Theorem 1]. See also [DK] Chapter 6] for a discussion on the role of this result in density Ramsey theory.

Remark 4.4. Lemma 4.2 can be used to relate the numerical invariants \( \text{PHJ}(k, \varepsilon) \) and \( \text{DHJ}(k, \delta) \) associated with the two versions of the density Hales–Jewett theorem. Indeed, notice that for every integer \( k \geq 2 \) and every \( 0 < \theta < \varepsilon \leq 1 \) we have

\[
(4.3) \quad \text{PHJ}(k, \varepsilon) \leq \text{DHJ}(k, \varepsilon) \leq (\varepsilon - \theta)^{-1} \cdot \text{PHJ}(k, \theta) \cdot k^{\text{PHJ}(k, \theta)}.
\]

4.1.2. Step 2: obtaining correlation with an insensitive set. We start by introducing the combinatorial analogue of the notion of an insensitive set. We note that this combinatorial analogue in fact predates Definition 1.5.

Definition 4.5. Let \( A \) be a finite set with \( |A| \geq 2 \), let \( n \) be a positive integer, and let \( \alpha, \beta \in A \) with \( \alpha \neq \beta \).

1. We say that a subset \( E \) of \( A^n \) is \( (\alpha, \beta) \)-insensitive if for every \( s, t \in A^n \) which are \( (\alpha, \beta) \)-equivalent we have that \( t \in E \) if and only if \( s \in E \).

2. We say that a subset \( E \) of an \( n \)-dimensional combinatorial space \( V \) of \( A^N \) is \( (\alpha, \beta) \)-insensitive in \( V \) if \( I_v^{-1}(E) \) is \( (\alpha, \beta) \)-insensitive, where \( I_v : A^n \to V \) denotes the canonical isomorphism associated with \( V \).

The following lemma is the second step of the proof of Theorem 4.1. It is precisely in the proof of this step that Theorem 1.6 is applied.
Lemma 4.6. Let $m \geq k \geq 2$ be integers, and let $0 < \delta \leq 1$. Set
\begin{equation}
N = \text{GRL}(k + 1, m + 1, \lceil 2(k + 1)4^k\delta^{-k} \rceil_{(2^{k+1})^{-1}})
\end{equation}
and let $n$ be a positive integer such that
\begin{equation}
n \geq \frac{2(k + 1)4^k}{\delta^{k+1}} (k + 1)^N N.
\end{equation}
Let $A$ be a set with $|A| = k + 1$, and let $D \subseteq A^n$ with $|D| \geq \delta|A^n|$. Then, either
(i) $D$ contains a combinatorial line of $A^n$, or
(ii) there exist $\beta \in A$, an $m$-dimensional combinatorial subspace $V$ of $A^n$ and a subset $S$ of $V$ with the following properties.
(a) $S = \bigcap_{\alpha \in A \setminus \beta} \mathcal{E}^{\alpha}$ where for every $\alpha \in A \setminus \{\beta\}$ the set $\mathcal{E}^{\alpha}$ is $(\alpha, \beta)$-insensitive in $V$.
(b) We have
\begin{equation}
|S| \geq \frac{\delta^{2k+1}}{(k + 1)^2 4^k + 2} \quad \text{and} \quad |D \cap S| \geq \delta + \frac{\delta^{k+1}}{(k + 1) 4^k + 1}.
\end{equation}

Remark 4.7. Lemma 4.6 improves upon two important quantitative aspects of what was known before. Firstly, by Proposition 2.2, the threshold on the dimension $n$ appearing in (4.5) is bounded by a primitive recursive function which belongs to the class $\mathcal{E}^5$ of Grzegorczyk’s hierarchy; in particular, it is independent of the numbers DHJ($k, \delta$). Secondly, the increment of the density of the set $D$ obtained in the second part of (4.6) depends polynomially on $\delta$; in order to appreciate this particular improvement we recall that all previous proofs yield a density increment which has an Ackermann-type dependence with respect to $k$. We also note that this quantity controls the number of iterations needed to be performed in order to prove Theorem 4.1 and as such it has significant impact on the behavior of the density Hales–Jewett numbers.

Proof of Lemma 4.6. We set $\eta := \frac{\delta^{k+1}}{2(k + 1)^2 4^k}$. By Lemma 4.2 and (4.5), there exist $\ell \in \{m, \ldots, n-1\}$ and an $N$-dimensional combinatorial subspace $V_1$ of $A^n$ such that for every $t \in V_1$ we have
\begin{equation}
\frac{|D_t|}{|A^{n-\ell}|} \geq \frac{|D|}{|A^n|} - \eta.
\end{equation}
We view the set $A^{n-\ell}$ as a discrete probability measure equipped with the uniform probability measure which we shall denote by $\mathbb{P}_1$. By Fact 3.1 and (4.4), there exists an $(m+1)$-dimensional combinatorial subspace $V_2$ of $V_1$ such that the process $\langle D_{t_{V_2(t)}} : t \in A^{m+1} \rangle$ is $\eta$-stationary; consequently, by Remark 3.3 and (4.7) and the fact that $|D| \geq \delta|A^n|$, there exists $\varepsilon \geq \delta$ such that $|\mathbb{P}_1(D_t) - \varepsilon| \leq \eta$ for every $t \in V_2$.

Now assume that part (i) does not hold true, that is, the set $D$ contains no combinatorial line of $A^n$. This in turn implies that $\bigcap_{\ell \in L} D_t = \emptyset$ for every combinatorial line $L$ of $V_2$; in particular, $\varepsilon \leq 1 - \frac{1}{2^{k+1}}$. Next, set $\sigma := \frac{\delta^{k+1}}{2(k + 1)^2 4^k}$ and notice that $\eta \leq \sigma/4^k$. Thus, by Theorem 1.6 there exist a nonempty subset $\Gamma$ of $A$, $\beta \in A \setminus \Gamma$
and a stochastic process \((S_{E_{t}}(t) : t \in A^{m+1})\) consisting of subsets of \(A^{n-\ell}\) such that the following are satisfied.

(a) For every \(t \in V_{2}\) we have \(S_{t} = \bigcap_{\alpha \in \Gamma} E_{t}^{\alpha}\) where for every \(\alpha \in \Gamma\) the stochastic process \((E_{t}^{\alpha} : t \in A^{m+1})\) is \((\alpha, \beta)\)-insensitive.

(b) For every \(t \in V_{2}\) such that \(I_{v_{2}}^{-1}(t)\) contains \(\beta\) we have

\[
P(S_{t}) \geq \frac{\varepsilon^k}{4(k + 1)} \quad \text{and} \quad P(D_{t} \ | \ S_{t}) \geq \varepsilon + \frac{\sigma}{4k}.
\]

By setting \(E_{t}^{\alpha} = A^{n-\ell}\) for every \(t \in V_{2}\) and every \(\alpha \in A \setminus (\Gamma \cup \{\beta\})\), we may assume that \(\Gamma = A \setminus \{\beta\}\). Next, let \(V_{3}\) denote the set of all \(t \in V_{2}\) such that \(I_{v_{2}}^{-1}(t)\) starts with \(\beta\), and notice that \(V_{3}\) is an \(m\)-dimensional combinatorial subspace of \(V_{2}\). Also observe that property (a) above and (4.8) hold true for every \(t \in V_{3}\).

With the process \((S_{t} : t \in V_{3})\) at our disposal the rest of the proof follows by a double counting argument and an application of the first moment method. Indeed, let \(P_{2}\) and \(P_{3}\) denote the uniform probability measures on \(V_{3}\) and \(V_{3} \times A^{n-\ell}\) respectively. Set \(S := \bigcup_{t \in V_{3}} \{ t \times s \in S \} \subseteq V_{3} \times A^{n-\ell}\) and notice that, by (4.8),

\[(4.9) \quad P_{3}(S) \geq \frac{\varepsilon^k}{4(k + 1)} \quad \text{and} \quad P_{3}(D_{t} \ | \ S) \geq \varepsilon + \frac{\sigma}{4k}.
\]

For every \(s \in A^{n-\ell}\) let \(S^{s} = \{ t \in V_{3} : t \cap s \in S \}\) and \(D^{s} = \{ t \in V_{3} : t \cap s \in D \}\) denote the sections of \(S\) and \(D\) at \(s\) respectively, and set

\[
B := \left\{ s \in A^{n-\ell} : P_{2}(S^{s}) \leq \frac{\varepsilon^k \sigma}{2(k + 1)4^k + 1} \right\} \quad \text{and} \quad C := \bigcup_{s \in B} S^{s} \times \{ s \} \subseteq S.
\]

Noticing that \(P_{3}(C) \leq (\varepsilon^k \sigma)/(2(k + 1)4^k + 1)\), by (4.9), we obtain that

\[(4.10) \quad P_{3}(C \ | \ S) \leq \frac{\sigma}{2 \cdot 4^k}.
\]

We thus have

\[
P_{3}(D \ | \ S \ \setminus \ C) = \frac{P_{3}(D \cap (S \ \setminus \ C))}{P_{3}(S \ \setminus \ C)} = \frac{P_{3}(D \cap (S \ \setminus \ C))}{P_{3}(S)} \geq P_{3}(D \ | \ S) - P_{3}(C \ | \ S) \geq \varepsilon + \frac{\sigma}{2 \cdot 4^k}.
\]

Since

\[
P_{3}(D \ | \ S \ \setminus \ C) = \sum_{s \in A^{n-\ell} \ \setminus \ B} P_{2}(D^{s} \ | \ S^{s}) \cdot P_{3}(S^{s} \times \{ s \} \ | \ S \ \setminus \ C)
\]

and \(\sum_{s \in A^{n-\ell} \ \setminus \ B} P_{3}(S^{s} \times \{ s \} \ | \ S \ \setminus \ C) = 1\), there exists \(s \in A^{n-\ell} \ \setminus \ B\) such that \(P_{2}(D^{s} \ | \ S^{s}) \geq \varepsilon + \sigma/(2 \cdot 4^k)\). We set \(V := V_{3} \times \{ s \}\), \(S := S \cap V\) and \(E^{\alpha} := \left( \bigcup_{t \in V_{3}} \{ t \times E_{t}^{\alpha} \} \right) \cap V\) for every \(\alpha \in A \setminus \{\beta\}\).

It is easy to see that with these choices the second part of the lemma is satisfied. The proof is completed. \(\square\)
4.1.3. Step 3: partitioning the insensitive set into combinatorial subspaces. The following lemma, which is proved in \cite[Lemma 8.2]{P2}, is the last step of the proof of Theorem 4.1.

**Lemma 4.8.** Let $k \geq 2$ be an integer, and assume that for every $0 < \delta \leq 1$ the number $\text{DHJ}(k, \delta)$ has been defined.

Then for every positive integer $m$ and every $0 < \eta \leq 1$ there exists a positive integer $\text{Til}(k, m, \eta)$—which depends on the numbers $\text{DHJ}(k, \delta)$—satisfying the following property. Let $A$ be a set with $|A| = k + 1$, let $n \geq \text{Til}(k, m, \eta)$ be an integer, and let $\beta \in A$. Also let $V$ be an $n$-dimensional combinatorial subspace of $A^{<\mathbb{N}}$ and let $S \subseteq V$ which is of the form $S = \bigcap_{\alpha \in A \setminus \{\beta\}} E^\alpha$ where $E^\alpha$ is $(\alpha, \beta)$-insensitive in $V$ for every $\alpha \in A \setminus \{\beta\}$. Then there exists a (possibly empty) collection $W$ of pairwise disjoint $m$-dimensional combinatorial subspaces of $V$ with $\bigcup W \subseteq S$ and such that $|S \setminus \bigcup W| \leq \eta|V|$.

Although the proof of Lemma 4.8 given in \cite{P2} is quite natural, unfortunately it leads to a very bad dependence of the numbers $\text{Til}(k, m, \eta)$ on the numbers $\text{DHJ}(k, \delta)$—see, e.g., \cite[Section 9]{P2} for a discussion on this issue.

**4.1.3. Completion of the proof of Theorem 4.1.** Given Lemmas 4.6 and 4.8, the proof of Theorem 4.1 follows easily by induction on $k$. (The base case “$k = 2$” is a consequence of the classical Sperner theorem \cite{Sp}. See, e.g., \cite[Chapter 8]{DK} or \cite{P2} for detailed expositions.

4.2. Comments. As alluded to earlier, Lemma 4.6 is a step towards obtaining primitive recursive bounds for the numbers $\text{DHJ}(k, \delta)$. It is clear that what is missing at this point is a quantitatively not wasteful proof of Lemma 4.8 (or a related variant). Although this will certainly require new ideas, it is likely that this program will eventually lead to primitive recursive bounds for the numbers $\text{DHJ}(k, \delta)$ belonging to the class $\mathcal{E}^7$ of Grzegorczyk’s hierarchy or slightly higher.

A disadvantage of this approach is that it relies on an analysis which is “local” in nature because we assume stationarity. It would be much more desirable if we had a “global” structure theorem. Formulating and proving a “global” theorem with quantitative aspects comparable to that of Theorem 1.6 might lead to upper bounds for the numbers $\text{DHJ}(k, \delta)$ which are of tower-type; note that this would also improve the longstanding upper bounds for the coloring version of Hales–Jewett theorem obtained by Shelah \cite{Sh}.

However, even tower-type upper bounds are rather unlikely to be anywhere close to optimal. Indeed, the best known lower bounds for the numbers $\text{DHJ}(k, \delta)$ are merely quasi-polynomial with respect to $\delta^{-1}$ (see \cite[Theorem 1.3]{P1}).
5. THE TYPE OF A SUBSET OF A DISCRETE HYPERCUBE

This is the first section of the second part of this paper which is devoted to the study of correlations of stochastic processes over arbitrary nonempty subsets of discrete hypercubes. As we have pointed out in the introduction, the analysis of these correlations relies, in a essentially way, on the notion of the \textit{type} of a nonempty subset of $A^n$. This Ramsey-theoretic invariant was introduced in \cite{DKT2}, though it can be traced\footnote{More precisely, the results in \cite{FK1} concern colorings of variable words—this is a similar, but not identical, setting.} in \cite{FK1}. We point out that for technical reasons (that will become transparent in Sections 6, 7 and 8), we will work with nonempty tuples of distinct elements of hypercubes instead of nonempty finite sets. This is an equivalent framework, but it does have some impact on our exposition when compared with that in \cite{DKT2}. With this machinery at our disposal, it is straightforward to extend the notions of stationarity, pseudorandomness, supercorrelation and subcorrelation introduced in Definitions 1.4 and 3.5 respectively; these extensions are presented in Subsection 5.4.

5.1. The type of a nonempty tuple. Let $A$ be a finite set with $|A| \geq 2$, and let $n, p$ be positive integers with $p \leq |A|^n$. Let $t = (t_1, \ldots, t_p)$ be a nonempty tuple (a nonempty finite sequence) of distinct elements of $A^n$.

5.1.1. If $p = 1$, then we define the type $\tau(t)$ of $t$ to be the empty sequence.

5.1.2. If $p \geq 2$, then we define $\tau(t)$ as follows. Let $R = (r_{ij}) \in A^{n \times p}$ denote the $n \times p$ matrix whose $(i,j)$-th entry $r_{ij}$ is the $i$-th coordinate of $t_j$. (More precisely, writing $t_j = (t_{i,j}, \ldots, t_{n,j})$ for every $j \in [p]$, we have $r_{ij} = t_{i,j}$.) Next, let $E$ denote the matrix which is obtained by first erasing all rows of $R$ with constant entries, and then shrinking all consecutive appearances of identical rows to single rows; note that $E$ is nonempty since $p \geq 2$. Let $m$ denote the numbers of rows of $E$, and let $s_1, \ldots, s_p$ denote its columns (in particular, we have that $s_j \in A^m$ for every $j \in [p]$). We define the type $\tau(t)$ of $t$ by the rule

\begin{equation}
\tau(t) = (s_1, \ldots, s_p)
\end{equation}

and we call the positive integer $m$ as the \textit{dimension} of $\tau(t)$. (Thus, $\tau(t)$ is a $p$-tuple of distinct elements of $A^m$.)

\textit{Example 5.1.} Let $A = [4]$, $n = 5$, $p = 5$, and

$$t = ((2, 1, 3, 2, 3), (3, 1, 4, 2, 4), (4, 1, 3, 2, 3), (3, 1, 4, 2, 4), (4, 1, 2, 2, 2)).$$
Then we have
\[
R = \begin{bmatrix} 2 & 3 & 4 & 3 & 4 \\ 1 & 1 & 1 & 1 & 1 \\ 3 & 4 & 3 & 4 & 2 \\ 2 & 2 & 2 & 2 & 2 \\ 3 & 4 & 3 & 4 & 2 \end{bmatrix}
\text{ and } E = \begin{bmatrix} 2 & 3 & 4 & 3 & 4 \\ 3 & 4 & 3 & 4 & 2 \end{bmatrix}
\]
and, consequently, \(m = 2\) and \(\tau(t) = ((2, 3), (3, 4), (4, 3), (3, 4), (4, 2))\).

**Example 5.2.** Let \(A\) be a finite set with \(|A| \geq 2\), let \(\Gamma \subseteq A\) be nonempty, set \(p := |\Gamma|\), and let \((\gamma_1, \ldots, \gamma_p)\) be an enumeration of the set \(\Gamma\). Also let \(n\) be an arbitrary positive integer. Then for every variable word \(v\) over \(A\) of length \(n\) we have \(\tau((v(\gamma_1), \ldots, v(\gamma_p))) = (\gamma_1, \ldots, \gamma_p)\).

We isolate, for future use, two basic properties of types which are both straightforward consequences of the definition. The first property shows that the type is an isomorphic invariant.

**Fact 5.3.** Let \(A\) be a finite set with \(|A| \geq 2\), let \(n, p\) be positive integers with \(2 \leq p \leq |A|^n\), and let \((t_1, \ldots, t_p)\) be a nonempty tuple of distinct elements of \(A^n\). Then for every \(n\)-dimensional combinatorial space \(V\) of \(A^{\subseteq N}\) we have
\[
\tau((t_1, \ldots, t_p)) = \tau((I_V(t_1), \ldots, I_V(t_p)))
\]
where \(I_V : A^n \to V\) denotes the canonical isomorphism associated with \(V\).

The second property is the permutation invariance of types.

**Fact 5.4.** Let \(A, n\) and \(p\) be as in Fact 5.3. Let \((t_1, \ldots, t_p)\) be a nonempty tuple of distinct elements of \(A^n\) and write \(\tau((t_1, \ldots, t_p)) = (s_1, \ldots, s_p)\). Then for every permutation \(\pi \in S_p\) we have \(\tau((t_{\pi(1)}, \ldots, t_{\pi(p)})) = (s_{\pi(1)}, \ldots, s_{\pi(p)})\).

5.2. The type of a nonempty finite set. Let \(A\) be a finite set with \(|A| \geq 2\), and let \(n\) be a positive integer. Let \(G \subseteq A^n\) be nonempty. Set \(p := |G|\) and fix an enumeration \((t_1, \ldots, t_p)\) of \(G\). If \(p = 1\) (that is, if \(G\) is a singleton), then we define the type \(\tau(G)\) of \(G\) to be the empty set. Otherwise, if \(p \geq 2\), then write \(\tau((t_1, \ldots, t_p)) = (s_1, \ldots, s_p)\) and define the type \(\tau(G)\) of \(G\) by setting
\[
\tau(G) = \{s_1, \ldots, s_p\}.
\]
Note that, by Fact 5.4, \(\tau(G)\) is well-defined and independent of the enumeration of \(G\), and observe that \(\tau(G)\) is a subset of \(A^m\) of cardinality \(|G|\) where \(m\) denotes the dimension of \(\tau((t_1, \ldots, t_p))\). By slightly abusing the previous terminology, we will call this positive integer \(m\) as the *dimension* of \(\tau(G)\). (Note that the dimension of \(\tau(G)\) controls its cardinality; specifically, we have \(|\tau(G)| \leq |A|^m\).) We set
\[
\text{Type}(A) := \{\tau(G) : G \text{ is a nonempty subset of } A^n \text{ for some integer } n \geq 1\}
\]
and we call an element of Type$(A)$ as a type over $A$. We also observe the following analogue of Fact 5.3. (As before, the proof is straightforward.)

**Fact 5.5.** Let $A, n$ and $V$ be as in Fact 5.3. Then for every nonempty $G \subseteq A^n$ we have $\tau(G) = \tau(1_V(G))$.

5.3. **Types and the Ramsey property.** The most important property of types is that they can be used in order to classify all partition regular families of subsets of discrete hypercubes. To motivate this classification, we start by observing that there is no analogue of Ramsey’s classical theorem for colorings of subsets of combinatorial spaces of a fixed cardinality. Indeed, let $A$ be a finite set with $|A| \geq 2$, and let $d, \ell \in \mathbb{N}$ with $|A|^d \geq \ell \geq 2$. Also let $V$ be a combinatorial space of $A^{|A|^d}$ of dimension at least $d + 1$, and define a coloring $c$ of the set \{ $G \subseteq V : |G| = \ell$ \} as follows. Let $G \subseteq V$ with $|G| = \ell$, and set $c(G) = \tau(G)$ if the dimension of the type of $G$ is at most $d$; otherwise set $c(G) = 0$. Regardless of how large the dimension of $V$ is, using Fact 5.3 it is easy to see that for every $(d + 1)$-dimensional combinatorial subspace $W$ of $V$ the set \{ $G \subseteq W : |G| = \ell$ \} is not monochromatic.

However, colorings which depend on the type are the only obstacles to the Ramsey property. Specifically, we have the following theorem whose proof can be found in [DK, Theorem 5.5] and which relies on the Graham–Rothschild theorem [GR].

**Theorem 5.6.** For every triple $k, m, r$ of positive integers with $k \geq 2$ there exists a positive integer $N$ with the following property. For every integer $n \geq N$, every set $A$ with $|A| = k$ and every $r$-coloring of the powerset of $A^n$ there exists an $m$-dimensional combinatorial subspace $V$ of $A^n$ such that every pair of nonempty subsets of $V$ with the same type is monochromatic. The least positive integer $N$ with this property is denoted by RamSp$(k, m, r)$.

Moreover, the numbers RamSp$(k, m, r)$ are upper bounded by a primitive recursive function belonging to the class $E^6$ of Grzegorczyk’s hierarchy.

5.4. **Stochastic processes and types: stationarity, pseudorandomness, supercorrelation, subcorrelation.** Our next goal is to extend Definitions 1.4 and 3.3. We begin by generalizing the notion of stationarity.

**Definition 5.7.** Let $A$ be a finite set with $|A| \geq 2$, let $n$ be a positive integer, let $\eta > 0$, and let $(D_t : t \in A^n)$ be a stochastic process in a probability space $(\Omega, \mathcal{F}, P)$. We say that $(D_t : t \in A^n)$ is $\eta$-stationary if for every pair of nonempty sets $G_1, G_2 \subseteq A^n$ with $\tau(G_1) = \tau(G_2)$ we have

$$(5.4) \quad \left| P\left( \bigcap_{t \in G_1} D_t \right) - P\left( \bigcap_{t \in G_2} D_t \right) \right| \leq \eta.$$  
(In particular, by Example 5.2, if a process $(D_t : t \in A^n)$ is $\eta$-stationary, then it is also $\eta$-stationary with respect to combinatorial lines.)

The following fact, which extends Fact 3.1 is an immediate consequence of Theorem 5.6.
Fact 5.8. Let $k \geq 2$ be an integer, and let $A$ be a set with $|A| = k$. Also let $0 < \eta \leq 1$, and let $n, m$ be positive integers such that

$$n \geq \text{RamSp}(k, m, [1/\eta]^{2k-1}).$$

Then for every stochastic process $\langle D_t : t \in A^n \rangle$ in a probability space $(\Omega, \mathcal{F}, \mathbb{P})$ there exists an $m$-dimensional combinatorial subspace $V$ of $A^n$ such that the process $\langle D_{1V(s)} : s \in A^m \rangle$ (that is, the restriction of $\langle D_t : t \in A^n \rangle$ on $V$) is $\eta$-stationary.

We also have the following analogue of Lemma 3.2 whose proof is identical to that of Lemma 3.2.

Lemma 5.9. Let $A, n, \eta$ and $\langle D_t : t \in A^n \rangle$ be as in Definition 5.7 Then the following are satisfied.

(i) For every $t_1, t_2 \in A^n$ we have $|\mathbb{P}(D_{t_1}) - \mathbb{P}(D_{t_2})| \leq \eta$. Thus, for every $t \in A^n$ we have $|\mathbb{P}(D_t) - \varepsilon| \leq \eta$ where $\varepsilon := \max \{ \max \{ \mathbb{P}(D_t) : t \in A^n \} : \eta \} > 0$.

(ii) Let $m \in [n]$, and let $\tau \in \text{Type}(A)$ be a type over $A$ of dimension $m$ and with $|\tau| \geq 2$. Then for every $Q \subseteq \tau$ and every pair $V_1, V_2$ of $m$-dimensional combinatorial subspaces of $A^n$ we have

$$\left| \mathbb{P}\left( \bigcap_{t \in V_1(Q)} D_t \cap \bigcap_{t \in (\tau \setminus Q)} D_t^c \right) - \mathbb{P}\left( \bigcap_{t \in V_2(Q)} D_t \cap \bigcap_{t \in (\tau \setminus Q)} D_t^c \right) \right| \leq 2^{1|Q|} \eta.$$

We proceed by generalizing the notions of pseudorandomness, supercorrelation and subcorrelation introduced in Definition 3.5.

Definition 5.10. Let $A$ be a finite set with $|A| \geq 2$, let $n$ be a positive integer, let $0 < \eta, \varepsilon \leq 1$, and let $\langle D_t : t \in A^n \rangle$ be an $\eta$-stationary process in a probability space $(\Omega, \mathcal{F}, \mathbb{P})$ such that $|\mathbb{P}(D_t) - \varepsilon| \leq \eta$ for every $t \in A^n$. Also let $\tau \in \text{Type}(A)$ be a type over $A$ of dimension at most $n$, and let $\theta \geq 0$.

1. (Pseudorandomness) We say that $\langle D_t : t \in A^n \rangle$ is $(\tau, \theta)$-pseudorandom if $|\mathbb{P}(\bigcap_{t \in G} D_t) - \varepsilon| \leq \theta$ for every $G \subseteq A^n$ with $\tau(G) = \tau$.
2. (Supercorrelation) We say that $\langle D_t : t \in A^n \rangle$ is $(\tau, \theta)$-supercorrelated if $\mathbb{P}(\bigcap_{t \in G} D_t) \geq \varepsilon + \theta$ for every $G \subseteq A^n$ with $\tau(G) = \tau$.
3. (Subcorrelation) We say that $\langle D_t : t \in A^n \rangle$ is $(\tau, \theta)$-subcorrelated if $\mathbb{P}(\bigcap_{t \in G} D_t) \leq \varepsilon - \theta$ for every $G \subseteq A^n$ with $\tau(G) = \tau$.

We close this section with the following analogue of Fact 3.6. (Its simple proof is left to the interested reader.)

Fact 5.11. Let $A, n, \eta, \varepsilon$ and $\langle D_t : t \in A^n \rangle$ be as in Definition 5.10. Also let $\tau \in \text{Type}(A)$ be a type over $A$ of dimension at most $n$, and let $\theta \geq \eta$. Then one of the following holds true.

(i) The process $\langle D_t : t \in A^n \rangle$ is $(\tau, \theta)$-pseudorandom.

(ii) The process $\langle D_t : t \in A^n \rangle$ is $(\tau, \theta - \eta)$-supercorrelated.

(iii) The process $\langle D_t : t \in A^n \rangle$ is $(\tau, \theta - \eta)$-subcorrelated.
6. The separation index

This section, like Section 5, also contains preparatory material which is needed for the analysis of arbitrary correlations of stationary stochastic processes. Our aim is to define another isomorphic invariant of nonempty subsets of discrete hypercubes—the separation index—which is coarser than the type, and measures how “well-distributed” a subset is. Specifically, we have the following definition.

**Definition 6.1.** Let $A$ be a finite set with $|A| \geq 2$, and let $n$ be a positive integer.

1. Let $\mathbf{t} = (t_1, \ldots, t_p)$ be a nonempty tuple of distinct elements of $A^n$, and let $\ell$ be a positive integer. We say that $\mathbf{t}$ is $\ell$-separated if for every $j \in [p]$ with $j \geq 2$ there exists $I \subseteq [n]$ (depending, possibly, on $j$) with $|I| = \ell$ and satisfying the following property: for every $q \in \{1, \ldots, j - 1\}$ there exists $i \in I$ such that $t_j(i) \neq t_q(i)$. (Namely, the $i$-th coordinate $t_j(i)$ of $t_j$ is different from the $i$-th coordinate $t_q(i)$ of $t_q$.) We define the separation index $s(\mathbf{t})$ of $\mathbf{t}$ to be the least positive integer $\ell$ such that $\mathbf{t}$ is $\ell$-separated.

2. Let $G \subseteq A^n$ be nonempty, and set $p := |G|$. We define the separation index $s(G)$ of $G$ by the rule

$$s(G) := \min\{s(\mathbf{t}) : \mathbf{t} = (t_1, \ldots, t_p) \text{ is an enumeration of } G\},$$

and we say that $G$ is $\ell$-separated if $s(G) = \ell$.

**Remark 6.2.** Note that the separation index of a nonempty finite set may be strictly smaller than the separation index of one of its enumerations. For instance, let $G = \{(0,0), (1,0), (0,1)\} \subseteq \{0,1\}^2$ and $\mathbf{t} = ((1,0), (0,1), (0,0))$. Then we have $s(\mathbf{t}) = 2$, but $s(G) = 1$ as witnessed by the tuple $\mathbf{s} = ((0,0), (1,0), (0,1))$.

In the following fact we state two basic properties of the separation index which were mentioned above, namely that it is preserved under canonical isomorphisms and that it is coarser than the type. The proof follows from the relevant definitions and is left to the reader.

**Fact 6.3.** Let $A$ be a finite set with $|A| \geq 2$, let $n$ be a positive integer, and let $V$ be an $n$-dimensional combinatorial space of $A^{<\mathbb{N}}$. If $(t_1, \ldots, t_p)$ is a nonempty tuple of distinct elements of $A^n$, then $s((t_1, \ldots, t_p)) = s((1_V(t_1), \ldots, 1_V(t_p)))$. Respectively, if $G \subseteq A^n$ is nonempty, then $s(G) = s(1_V(G))$; consequently, if $H \subseteq A^l$ for some positive integer $l$ with $\tau(H) = \tau(G)$, then $s(H) = s(G)$.

We proceed by determining the separation index of some concrete examples of sets which are important from a combinatorial perspective.

**Example 6.4 (Combinatorial lines).** Let $A$ and $n$ be as in Definition 6.1. Let $\Gamma \subseteq A$ be nonempty, and set $p := |\Gamma|$. Also let $v$ be a variable word over $A$ of length $n$. Then, by Fact 6.3, for every enumeration $(\gamma_1, \ldots, \gamma_p)$ of $\Gamma$ we have $s((v(\gamma_1), \ldots, v(\gamma_p))) = s((\gamma_1, \ldots, \gamma_p)) = 1$. In particular, every combinatorial line $L$ of $A^n$ is 1-separated.
Example 6.5 (Shelah lines). As above, let $A$ be a finite set with $|A| \geq 2$. For every $\alpha \in A$ and every positive integer $m$ let $\alpha^m = (\alpha, \ldots, \alpha)$ denote the sequence of length $m$ taking the constant value $\alpha$; also let $\alpha^0$ denote the empty sequence.

Now let $n$ be a positive integer, let $\alpha, \beta \in A$ with $\alpha \neq \beta$, and define the Shelah line\footnote{These sets play a crucial role in Shelah’s proof of the Hales–Jewett theorem.} with parameters $\alpha, \beta$ by rule

\begin{equation}
S = \{\alpha^{n-m} \beta^m : m \in \{0, \ldots, n\}\} \subseteq A^n.
\end{equation}

Clearly, we have $|S| = n + 1$, and it is easy to see that the set $S$ is 1-separated.

Example 6.5 implies, in particular, that there exist 1-separated sets of arbitrarily large cardinality. More generally, we have the following lemma.

Lemma 6.6 (Random tuples of small size are 1-separated). Let $k, n, p$ be positive integers with $k \geq 2$ and $2 \leq p \leq n^k$. Let $A$ be a set with $|A| = k$, and let $\mathbb{P}$ denote the uniform probability measure on $(A^n)^p$. (That is, $(A^n)^p$ is the Cartesian product of $p$ many copies of $A^n$.) Then we have

\begin{equation}
\mathbb{P}(t \text{ is 1-separated}) \geq 1 - pe^{-n(k+1)^p}.
\end{equation}

In particular, if $p \leq \log(n)$, then $\mathbb{P}(t \text{ is 1-separated}) = 1 - o_{n \to \infty;k}(1)$.

Proof. Set $S := \{t \in (A^n)^p : t \text{ is 1-separated}\}$, and let $S^c$ denote the complement of $S$. Note that for every $i \in [n]$ and every $j \in [p]$ with $j \geq 2$ the set of all $t = (t_1, \ldots, t_p) \in (A^n)^p$ such that $t_j(i) \notin \{t_1(i), \ldots, t_{j-1}(i)\}$ has probability $e^{-nk} \geq (k^{-1})^p$. (Here, $t_j(i)$ denotes the $i$-th coordinate of $t_j$ for every $q \in [j].$)

Therefore, for every $j \in [p]$ with $j \geq 2$ the set of all $t = (t_1, \ldots, t_p) \in (A^n)^p$ such that $t_j$ fails to satisfy the condition of being 1-separated has probability at most $(1 - (\frac{k}{n})^p)^n$. Using the fact that $(1 - \frac{r}{n})^n \leq e^{-r}$ for every $r > 0$ and every positive integer $n$, we thus have

\begin{equation}
\mathbb{P}(S^c) \leq p\left(1 - \left(\frac{k-1}{k}\right)^p\right)^n \leq pe^{-n(k+1)^p}
\end{equation}

which is equivalent to (6.3).

Next assume that $p \leq \log(n)$. Since the function $f(x) = xe^{-nx}$ is increasing for every $r \in (0,1)$ and every positive integer $n$, by (6.4), we obtain that

$$
\mathbb{P}(S^c) \leq \log(n) e^{-n(k+1)^p\log(n)} = \log(n) e^{-n^{1+\log(k+1)}}.
$$

Therefore, $\mathbb{P}(S) = 1 - o_{n \to \infty;k}(1)$ as desired. \hfill \Box

The last example in this section provides us with a representative example of an $n$-separated set.

Example 6.7 (Combinatorial subspaces). Let $A$ and $n$ be as in Definition 6.1 and notice that for every nonempty $G \subseteq A^n$ we have $s(G) \leq n$. On the other hand, it is easy to verify that $s(A^n) = n$. Using this observation and Fact 6.3, we see that every $n$-dimensional combinatorial space of $A^{<\omega}$ is $n$-separated.
7. Correlations over 1-separated sets

7.1. The main result. We begin by introducing the analogue of insensitivity for processes indexed by combinatorial spaces.

**Definition 7.1.** Let $A, n, \alpha$ and $\beta$ be as in Definition 1.5, let $V$ be an $n$-dimensional combinatorial space of $A^{<N}$, and let $I_V : A^n \to V$ denote the canonical isomorphism associated with $V$. We say that a stochastic process $\langle D_t : t \in V \rangle$ in a probability space $(\Omega, \mathcal{F}, P)$ is $(\alpha, \beta)$-insensitive in $V$ if $\langle D_{I_V(s)} : s \in A^n \rangle$ is $(\alpha, \beta)$-insensitive in the sense of Definition 1.5. (That is, if $D_{I_V(s)} = D_{I_V(t)}$ for every $s, t \in A^n$ which are $(\alpha, \beta)$-equivalent.)

The main result of this section is the following extension of Theorem 1.6 which concerns correlations of stationary processes over 1-separated sets. (We recall that the notion of stationarity in this more general context is given in Definition 5.7.)

**Theorem 7.2.** Let $k, \kappa, m$ be positive integers with $k, \kappa \geq 2$ and $\kappa \leq km$, and let $\varepsilon, \sigma, \eta > 0$ such that

\[
\varepsilon \leq 1 - \frac{1}{2\kappa}, \quad \sigma \leq \frac{\varepsilon^{\kappa-1}}{2\kappa} \quad \text{and} \quad \eta \leq \frac{\sigma}{4\kappa-1}.
\]

Also let $A$ be a set with $|A| = k$, let $n > m$ be an integer, and let $\langle D_t : t \in A^n \rangle$ be an $\eta$-stationary process in a probability space $(\Omega, \mathcal{F}, P)$ such that $|P(D_t) - \varepsilon| \leq \eta$ for every $t \in A^n$. Then, either

(i) for every nonempty 1-separated $G \subseteq A^n$ with cardinality at most $\kappa$ and whose type $\tau(G)$ has dimension at most $m$ we have

\[
\left| P\left( \bigcap_{t \in G} D_t \right) - \varepsilon^{|G|} \right| \leq \sigma,
\]

(ii) or $\langle D_t : t \in A^n \rangle$ correlates with a “structured” stochastic process when restricted on a large subspace; precisely, there exist a combinatorial subspace $V$ of $A^n$ with $\dim(V) \geq n - m$, a nonempty subset $\Gamma$ of $A \setminus \Gamma$ and a stochastic process $\langle S_t : t \in V \rangle$ in $(\Omega, \mathcal{F}, P)$ with the following properties.

(a) For every $t \in V$ we have $S_t = \bigcap_{\alpha \in \Gamma} E^\alpha_t$ where for every $\alpha \in \Gamma$ the process $\langle E^\alpha_t : t \in V \rangle$ is $(\alpha, \beta)$-insensitive in $V$.

(b) For every $t \in V$ we have

\[
P(S_t) \geq \frac{\varepsilon^{\kappa-1}}{4\kappa} \quad \text{and} \quad P(D_t | S_t) \geq \varepsilon + \frac{\sigma}{4\kappa-1}.
\]

Theorem 7.2 shows that stationary processes which exhibit non-independent behavior over 1-separated sets are essentially characterized—in the strong quantitative sense described in (7.3)—by their correlation with insensitive processes. Note, however, that in contrast to Theorem 1.6, this correlation is “local” in nature, that is, we need to pass to a subspace in order to verify it. We present an example in Subsection 7.2 which elucidates the necessity of this restriction.
The proof of Theorem 7.2 is given in Subsection 7.5. It relies on the following analogue of Proposition 3.7, whose proof is given in Subsection 7.4. (The concepts of pseudorandomness, supercorrelation and subcorrelation which appear below are introduced in Definition 5.10.)

**Proposition 7.3.** Let $A$ be a finite set with $|A| \geq 2$, let $n, p$ be positive integers with $p + 1 \leq |A|^n$, let $0 < \eta, \varepsilon \leq 1$, and let $\langle D_t : t \in A^n \rangle$ be an $\eta$-stationary process in a probability space $(\Omega, \mathcal{F}, \mathbb{P})$ such that $|\mathbb{P}(D_t) - \varepsilon| \leq \eta$ for every $t \in A^n$.

Let $t = (t_1, \ldots, t_p)$ be an 1-separated tuple consisting of distinct elements of $A^n$, set $G := \{t_1, \ldots, t_p+1\}$ and $H := \{t_1, \ldots, t_p\}$, and let $d$ denote the dimension of $\tau(G)$. Finally, let $0 < \theta, \sigma \leq 1$, and assume that the process $\langle D_t : t \in A^n \rangle$ is $(\tau(H), \theta)$-pseudorandom. Then there exist an $(n-d)$-dimensional combinatorial subspace $V$ of $A^n$, a nonempty subset $\Gamma$ of $A$, $\beta \in A \setminus \Gamma$ and a process $\langle S_t : t \in V \rangle$ in $(\Omega, \mathcal{F}, \mathbb{P})$ with the following properties.

(i) For every $t \in V$ we have $S_t = \bigcap_{\alpha \in \Gamma} E_t^\alpha$ where for every $\alpha \in \Gamma$ the process $\langle E_t^\alpha : t \in V \rangle$ is $(\alpha, \beta)$-insensitive in $V$.

(ii) For every $t \in V$ we have $|\mathbb{P}(S_t) - \varepsilon^{n^d}| \leq \theta$.

(iii) If $\langle D_t : t \in A^n \rangle$ is $(\tau(G), \sigma)$-supercorrelated, then for every $t \in V$ we have

$$\mathbb{P}(D_t | S_t) \geq \varepsilon \left(1 + \frac{\sigma \varepsilon^{-1} - \theta}{\varepsilon^{n^d} + \theta}\right).$$

(iv) If $\langle D_t : t \in A^n \rangle$ is $(\tau(G), \sigma)$-subcorrelated, then for every $t \in V$ we have

$$\mathbb{P}(D_t | S_t) \leq \varepsilon \left(1 - \frac{\sigma \varepsilon^{-1} - \theta}{\varepsilon^{n^d} - \theta}\right).$$

7.2. Correlations over 1-separated sets: example. We are about to present an example of a process which exhibits non-independent behavior when we look at its correlations over 1-separated sets whose type is rather simple, but not quite similar to that of combinatorial lines. As in Example 1.3 for concreteness we will work with the set $A = \{1, 2, 3\}$ and the 1-separated type

$$\tau = \{(1, 2, 1), (2, 1, 2), (2, 2, 3)\} \in \text{Type}(\{1, 2, 3\}).$$

Let $n \geq 5$ be an integer, and let

$$\langle E_{y^z} : y \in \{1, 2, 3\}^3 \setminus \{(2, 2, 3)\} \text{ and } s \in \{1, 2\}^{n-3}\rangle$$

be a family of independent events in a probability space $(\Omega, \mathcal{F}, \mathbb{P})$ with equal probability $\varepsilon > 0$. We define $\langle D_t : t \in \{1, 2, 3\}^n \rangle$ by setting

(a) $D_{(2,2,3)^{z-2}} := E_{(1,2,1)^{z-2行列}} \cap E_{(2,1,2)^{z-2行列}}$ for every $z \in \{1, 2, 3\}^{n-3}$, and

(b) $D_{y^z} := E_{y^z\cap E_{(2,1,2)^{z-2行列}}}$ if $y \in \{1, 2, 3\}^3 \setminus \{(2,2,3)\}$ and $z \in \{1, 2, 3\}^{n-3}$.

Note the difference between the definition in (a) and the definition in Example 1.3 given $t \in A^n$, first we change a short initial segment of $t$ and then we “project” the rest of the sequence. This maneuver will be generalized in the next subsection.

The analysis of the correlations of $\langle D_t : t \in \{1, 2, 3\}^n \rangle$ is fairly straightforward. Specifically, notice that if $t = y^z \in \{1, 2, 3\}^n$ with $y \in \{1, 2, 3\}^3 \setminus \{(2,2,3)\}$ and
$s \in \{1, 2\}^{n-3}$, then we have $D_t = E_t$ and, consequently, $P(D_t) = \varepsilon$; otherwise, we have $P(D_t) = \varepsilon^2$. (Thus, $P(D_t) = \varepsilon^2$ for “almost every” $t \in \{1, 2, 3\}^n$.) Moreover, for every $w \in \{1, 2, 3\}^{n-4}$ set

$$G_w := \{(1, 2, 1)^{(3^-w^{1-1})}, (2, 1, 2)^{(3^-w^{3-2})}, (2, 2, 3)^{(3^-w)}\} \subseteq \{1, 2, 3\}^n$$

and observe that $\tau(G_w) = \tau$ and $P(\bigcap_{t \in G_w} D_t) = \varepsilon^4$ which deviates, of course, from the expected value $\varepsilon^6$.

Finally, note that $\langle D_t : t \in \{1, 2, 3\}^n \rangle$ cannot be written as the intersection of insensitive processes, but only barely so. Indeed, set

$$(7.7) \quad V := \{(2, 2, 3)^{-z} : z \in \{1, 2, 3\}^{n-3}\}$$

and observe that $V$ is an $(n - 3)$-dimensional combinatorial subspace of $\{1, 2, 3\}^n$.

Clearly, by (a) above, the restriction of $\langle D_t : t \in \{1, 2, 3\}^n \rangle$ on $V$ is the intersection of two processes which are $(1, 3)$- and $(2, 3)$-insensitive in $V$ respectively.

7.3. Definitions/Notation. Let $A, n$ and $p$ be as in Proposition 7.3. Let

$$(7.8) \quad t = (t_1, \ldots, t_{p+1})$$

be an 1-separated tuple consisting of distinct elements of $A^n$, let $\tau = \tau(t)$ denote its type, and let $d$ denote the dimension of $\tau$. We will define

- an $(n - d)$-dimensional combinatorial subspace $V$ of $A^n$,
- a nonempty subset $\Gamma$ of $A$,
- $\beta \in A \setminus \Gamma$,
- an integer $\iota \in [d]$, and
- for every $j \in [p]$ a map $T_j : V \rightarrow A^n$.

These data will be used in the proofs of Theorem 7.2 and Proposition 7.3—in fact, they constitute the combinatorial heart of the argument. We also note that $V, \Gamma, \beta, \iota$ and $\langle T_j : j \in [p] \rangle$ will essentially depend upon the type $\tau$ of $t$ and not on the tuple $t$ itself; however, it is technically easier to work with $t$.

7.3.1. Defining $\iota, \beta$ and $\Gamma$, and splitting the type $\tau$. We write $\tau = (s_1, \ldots, s_{p+1})$ where $s_j = (s_j(1), \ldots, s_j(d)) \in A^d$ for every $j \in [p+1]$. By Fact 6.3 and our assumption that $t$ is 1-separated, we see that $\tau$ is also 1-separated. Taking into account this remark, we define

$$(7.9) \quad \iota := \min \{i \in [d] : s_{p+1}(i) \neq s_j(i) \text{ for every } j \in [p]\}$$

and

$$(7.10) \quad \beta := s_{p+1}(\iota), \quad \beta_j := s_j(\iota) \text{ for every } j \in [p], \quad \Gamma := \{\beta_1, \ldots, \beta_p\}.$$  

(In particular, we have $\beta \notin \Gamma$; also note that $|\Gamma| \leq p$ since the elements $\beta_1, \ldots, \beta_p$ are not necessarily distinct.) Moreover, for every $j \in [p+1]$ set

$$(7.11) \quad x_j = (s_j(1), \ldots, s_j(\iota)) \quad \text{and} \quad y_j = (s_j(\iota + 1), \ldots, s_j(n))$$

with the convention that $y_j$ is the empty sequence if $\iota = d$; note that $s_j = x_j^{\iota} y_j$. 
7.3.2. Defining $V$ and the maps $\langle T_j : j \in [p] \rangle$. Next, set

$$(7.12) \quad V := \{x_{p+1} \sim z \sim y_{p+1} : z \in A^{n-d}\}$$

and observe that $V$ is an $(n-d)$-dimensional combinatorial subspace\footnote{Notice that the subspace $V$ is of very special form; in particular, the canonical isomorphism associated with $V$ is the map $A^{n-d} \ni z \mapsto x_{p+1} \sim z \sim y_{p+1} \in V$.} of $A^n$.

Finally, for every $j \in [p]$ we define $T_j : V \to A^n$ by the rule

$$(7.13) \quad T_j (x_{p+1} \sim z \sim y_{p+1}) = x_j \sim (z^{\beta \to \beta_j}) \sim y_j.$$
that is, \( \langle S_t : t \in V \rangle \) satisfies part (ii) of the theorem.

For part (iii) assume that \( \langle D_t : t \in A^n \rangle \) is \( (\tau(G), \sigma) \)-super correlated, and let \( t \in V \) be arbitrary. By (7.14) and the super correlation assumption, we have that \( P(D_t \cap S_t) \geq \varepsilon^p + \sigma \). On the other hand, by (7.16), we see that \( P(S_t) \geq \varepsilon^p - \theta \). Therefore, \( P(D_t | S_t) \geq \varepsilon \left( 1 + \frac{\sigma - 1 - \theta}{\varepsilon^p + \sigma} \right) \) as desired.

Finally, for part (iv) assume that \( \langle D_t : t \in A^n \rangle \) is \( (\tau(G), \sigma) \)-subcorrelated. Fix \( t \in V \). Using again (7.14), the subcorrelation assumption and (7.16), we obtain that \( P(D_t \cap S_t) \leq \varepsilon^p - \sigma \) and \( P(S_t) \geq \varepsilon^p - \theta \) which implies that \( P(D_t | S_t) \leq \varepsilon \left( 1 - \frac{\sigma - 1 - \theta}{\varepsilon^p - \sigma} \right) \).

The proof is completed.

7.5. Proof of Theorem 7.2. It is similar to the proof of Theorem 1.6, with the main new ingredients being Proposition 7.3 and the material in Subsection 7.3. We shall describe in detail the necessary changes, as this proof will also serve as a model for the proof of Theorem 8.5 in Section 8.

Let \( (\theta_p)_{p=0}^\infty \) be the finite sequence defined in (3.6)—that is, \( \theta_0 = 0, \theta_1 = \eta \), and \( \theta_p = 4^{p-\kappa} \sigma \) if \( p \in \{2, \ldots, \kappa\} \)—and recall that \( (\theta_p)_{p=0}^\infty \) is increasing. Assume that part (i) of the theorem does not hold true, and fix an 1-separated set \( G \subseteq A^n \) of cardinality at most \( \kappa \) whose type \( \tau(G) \) has dimension at most \( m \) and such that:

(a) the process \( \langle D_t : t \in A^n \rangle \) is not \( (\tau(G), \theta_{|G|}) \)-pseudorandom, and (b) \( G \) has the minimal cardinal among all sets with these properties. (Note that \( |G| \geq 2 \).) Let \( t = (t_1, \ldots, t_{|G|}) \) be an enumeration of \( G \) such that the tuple \( t \) is 1-separated, let \( d \) denote the dimension of \( \tau(G) \), and set \( H := \{t_1, \ldots, t_{|G|-1}\} \) and \( p := |H| \); notice that \( 1 \leq p \leq \kappa - 1 \) and \( 1 \leq d \leq m \). Also observe that for every nonempty proper subset \( \Sigma \) of \( G \) the process \( \langle D_t : t \in A^n \rangle \) is \( (\tau(\Sigma), \theta_{|\Sigma|}) \)-pseudorandom.

We set \( \theta := \theta_p \) and \( \Theta := \theta_{p+1} \). Since \( \langle D_t : t \in A^n \rangle \) is not \( (\tau(G), \Theta) \)-pseudorandom, by Fact 5.11 we see that either

(A1) the process \( \langle D_t : t \in A^n \rangle \) is \( (\tau(G), \Theta - \eta) \)-super correlated,

(A2) or the process \( \langle D_t : t \in A^n \rangle \) is \( (\tau(G), \Theta - \eta) \)-subcorrelated.

If the first case holds true, then, arguing precisely as in the proof of Theorem 1.6 and using Proposition 7.3 instead of Proposition 3.7, it is easy to verify that part (ii) of the theorem is satisfied.

So assume that the process \( \langle D_t : t \in A^n \rangle \) is \( (\tau(G), \Theta - \eta) \)-subcorrelated, and let \( V \) and \( \langle T_j : j \in [p]\rangle \) be the combinatorial space and the maps obtained in Subsection 7.3 for the 1-separated tuple \( t \). For every \( t \in V \) we set

\[
S_t := \left( \bigcap_{j=1}^{p-1} D_{T_j(t)} \right) \cap D_{T_p(t)} = \left( \bigcap_{j=1}^{p-1} D_{T_j(t)} \right) \setminus \left( \bigcap_{j=1}^{p} D_{T_j(t)} \right).
\]

(Recall that, by convention, \( \bigcap_{j=1}^{p-1} D_{T_j(t)} = \Omega \) if \( p = 1 \).) Notice that, by part (ii) of Fact 7.4, the process \( \langle S_t : t \in V \rangle \) satisfies part (ii.a) of the theorem. Next, we set \( F := \{t_1, \ldots, t_{p-1}\} \) (observe that \( F \) may be empty). Since \( \langle D_t : t \in A^n \rangle \) is \( (\tau(F), \theta_{p-1}) \)-pseudorandom if \( p > 1 \) (if \( p = 1 \), then this is superfluous) and
Finally, by (7.1) and (7.18), we conclude that
\[(\tau(H), \theta)\text{-pseudorandom, by (7.15), for every } t \in V \text{ we have}
\]
\[
\|P(S_t) - \varepsilon^{p-1}(1 - \varepsilon)\| \leq \theta + \theta_{p-1}.
\]
Using (7.14) and the fact that the process \((D_t : t \in A^n)\) is \((\tau(G), \Theta - \eta)\text{-subcorrelated and } (\tau(F \cup \{t_{p+1}\}), \theta)\text{-pseudorandom, for every } t \in V \text{ we have}
\]
\[
\text{Moreover, by (7.1) and the definition of } (\theta_j)_{j=1}^\kappa, \text{ we see that } \theta + \theta_{p-1} \leq \varepsilon^p, \theta \leq \Theta/4 \text{ and } \theta_{p-1} + \eta \leq \Theta/4. \text{ Therefore, by (7.18) and (7.19), for every } t \in V
\]
\[
P(D_t \cap S_t) \geq (\varepsilon^p - \theta) - (\varepsilon^{p+1} - \Theta + \eta).
\]
\[
\text{Finally, by (7.1) and (7.18), we conclude that}
\]
\[
P(S_t) \geq \varepsilon^{p-1} - \varepsilon^p - \theta - \theta_{p-1} \geq \varepsilon^{\kappa-1}(1 - \varepsilon) - 2\theta \geq \frac{\varepsilon^{\kappa-1}}{2\kappa} - 2\theta \geq \frac{\varepsilon^{\kappa-1}}{4\kappa}
\]
for every \(t \in V\). The proof is completed.

8. Correlations over \(\ell\)-separated sets

8.1. Obstructions to independence: simplicial processes. We are about to begin our analysis of arbitrary correlations of stationary processes. As we have noted, the main—and perhaps the most interesting—difference lies in the fact that insensitive process are not enough to characterize non-independent behavior. Our goal in this subsection is to discuss this phenomenon and introduce the “structured” processes which appear in this more general context.

To this end, we need to define a “local” version of insensitivity. To motivate this “local” version, let \(A\) be a finite set with \(|A| \geq 2\), let \(n, \ell, r_1, \ldots, r_\ell\) be positive integers with \(n, \ell \geq 2\) and \(n = r_1 + \cdots + r_\ell\), and note that we may identify the hypercube \(A^n\) with the product \(A^{r_1} \times \cdots \times A^{r_\ell}\) via the map
\[
A^{r_1} \times \cdots \times A^{r_\ell} \ni (t_1, \ldots, t_\ell) \mapsto t_1 \land \cdots \land t_\ell \in A^n.
\]

Having in mind this identification, we may consider subsets of \(A^n\) which are insensitive only in one of the factors \(A^{r_1}, \ldots, A^{r_\ell}\). This is, essentially, the content of the following definition.

**Definition 8.1** (Local insensitivity). Let \(A\) be a finite set with \(|A| \geq 2\), let \(n\) be a positive integer, let \(\alpha, \beta \in A\) with \(\alpha \neq \beta\), and let \(I \subseteq [n]\) be nonempty. Also let \((\Omega, \mathcal{F}, P)\) be a probability space.

1. Let \(t, s \in A^n\) and write \(t = (t_1, \ldots, t_n)\) and \(s = (s_1, \ldots, s_n)\). We say that \(t, s\) are \((\alpha, \beta, I)\text{-equivalent if for every } i \in [n] \setminus I \text{ we have } t_i = s_i \text{ and, moreover, for every } i \in I \text{ and every } \gamma \in A \setminus \{\alpha, \beta\} \text{ we have } t_i = \gamma \text{ if and only if } s_i = \gamma\).
(2) We say that a process \( \langle E_t : t \in A^n \rangle \) in \( (\Omega, \mathcal{F}, \mathbb{P}) \) is \((\alpha, \beta, I)\)-insensitive if \( E_t = E_s \) for every \( t, s \in A^n \) which are \((\alpha, \beta, I)\)-equivalent.

(3) Let \( V \) be an \( n \)-dimensional combinatorial space of \( A^{< \mathbb{N}} \). We say that a process \( \langle E_t : t \in V \rangle \) in \( (\Omega, \mathcal{F}, \mathbb{P}) \) is \((\alpha, \beta, I)\)-insensitive in \( V \) provided that \( \langle E_{iV(t)} : t \in A^n \rangle \) is \((\alpha, \beta, I)\)-insensitive where \( I_V : A^n \to V \) denotes the canonical isomorphism associated with \( V \).

We proceed with the following example which shows the need to extend the notion of a “structured” process.

Example 8.2. As in Example 1.3, we will work with the set \( A = \{1, 2, 3\} \), and we will focus on correlations over 2-dimensional combinatorial spaces of \( \{1, 2, 3\}^{< \mathbb{N}} \). Notice that, by Fact 5.5 and Example 6.7, all 2-dimensional combinatorial spaces are 2-separated and are of type

\[ \tau = \{(1, 1), (1, 2), (1, 3), (2, 1), (2, 2), (2, 3), (3, 1), (3, 2), (3, 3)\} \in \text{Type}(\{1, 2, 3\}). \]

Now let \( n \) be an arbitrary positive integer, and fix a family

\[ \langle E_t^{-s} : t^s \in (\{1, 2\}^n \times \{1, 2, 3\}^n) \cup (\{1, 2, 3\}^n \times \{1, 2\}^n) \rangle \]

of independent events in a probability space \( (\Omega, \mathcal{F}, \mathbb{P}) \) with equal probability \( \varepsilon > 0 \).

We define a process \( \langle D_z : z \in \{1, 2, 3\}^{2n} \rangle \) by setting

\[
D_{t^{-s}} := S_{t^{-s}}^1 \cap S_{t^{-s}}^2
\]

where \( S_{t^{-s}}^1 := E_{t^3 \rightarrow t^1 \rightarrow s^3 \rightarrow 1} \cap E_{t^2 \rightarrow t^1 \rightarrow s^2 \rightarrow 2} \cap E_{t^3 \rightarrow t^3 \rightarrow 1} \cap E_{t^2 \rightarrow t^2 \rightarrow s^2 \rightarrow 2} \cap E_{t^1 \rightarrow t^2 \rightarrow s} \) and \( S_{t^{-s}}^2 := E_{t^2 \rightarrow t^3 \rightarrow 1} \cap E_{t^1 \rightarrow t^3 \rightarrow 2} \) for every \( t, s \in \{1, 2, 3\}^n \). Note that

\[
\mathbb{P}(D_{t^{-s}}) = \begin{cases} 
\varepsilon^8 & \text{if both } t, s \text{ contain 3,} \\
\varepsilon^3 & \text{if exactly one of } t, s \text{ contains 3,} \\
\varepsilon & \text{if both } t, s \text{ do not contain 3.}
\end{cases}
\]

Next, for every \( t, s \in \{1, 2, 3\}^n \) which contain 3 set

\[ G_{t,s} := \{t^3 \rightarrow 1 \rightarrow s^3 \rightarrow 1, t^1 \rightarrow 1 \rightarrow s^3 \rightarrow 2, t^3 \rightarrow 1 \rightarrow s^3 \rightarrow 1, t^1 \rightarrow 2 \rightarrow s^3 \rightarrow 1, t^1 \rightarrow 3 \rightarrow 2 \rightarrow s, t^3 \rightarrow 2 \rightarrow s, t^3 \rightarrow 3 \rightarrow 2, t^3 \rightarrow 3 \rightarrow s^3 \rightarrow 1 \}
\]

and observe that \( G_{t,s} \) is a 2-dimensional combinatorial subspace of \( \{1, 2, 3\}^{2n} \); also notice that \( D_{t^{-s}} = \bigcap_{z \in G_{t,s}} D_z \) and, therefore, \( \mathbb{P}(\bigcap_{z \in G_{t,s}} D_z) = \varepsilon^8 \) which deviates from the expected value \( \varepsilon^{24} \). In other words, the process \( \langle D_z : z \in \{1, 2, 3\}^{2n} \rangle \) exhibits non-independent behavior when we look at its correlations over 2-dimensional combinatorial subspaces.

Note, however, that \( \langle D_z : z \in \{1, 2, 3\}^{2n} \rangle \) cannot be written as the intersection of insensitive processes even if we restrict it on subspaces of very small dimension. (This is a consequence of the fact that the processes \( \langle S_{t^{-s}}^1 : t, s \in \{1, 2, 3\}^n \rangle \) and \( \langle S_{t^{-s}}^2 : t, s \in \{1, 2, 3\}^n \rangle \) depend non-trivially on the parameters \( s \) and \( t \) respectively.) Nevertheless, the process \( \langle D_z : z \in \{1, 2, 3\}^{2n} \rangle \) is not random at all: it is obtained
from \(\langle S^1_{t,s} : t,s \in \{1,2,3\}^n \rangle\) and \(\langle S^2_{t,s} : t,s \in \{1,2,3\}^n \rangle\) which are both the intersection of locally insensitive processes but for disjoint domains of insensitivity. This less restrictive form of structurability is abstracted in the following definition.

**Definition 8.3** (Simplicial processes). Let \(A\) be a finite set with \(|A| \geq 2\), let \(n, \ell\) be positive integers with \(n \geq \ell\), let \(r = (r_1, \ldots, r_\ell)\) be an \(\ell\)-tuple of positive integers such that \(n = r_1 + \cdots + r_\ell\), and let \(I^1_\ell, \ldots, I^r_\ell\) denote the unique successive intervals of \([n]\) such that \(|I^s_\ell| = r_\ell\) for every \(l \in [\ell]\). We say that a process \(\langle S_t : t \in A^n \rangle\) in a probability space \((\Omega, \mathcal{F}, \mathbb{P})\) is \((\ell, r)\)-simplicial if there exist

\[ \bullet \beta_1, \ldots, \beta_\ell \in A, \]
\[ \bullet \text{ for every } l \in [\ell] \text{ a nonempty subset } \Gamma_l \text{ of } A \setminus \{\beta_\ell\}, \] and
\[ \bullet \text{ for every } l \in [\ell] \text{ and } \alpha \in \Gamma_l \text{ an } (\alpha, \beta_1, I^s_\ell) \text{-insensitive process } \langle E^l_\alpha \rangle : t \in A^n, \]

such that for every \(t \in A^n\) we have

\[ S_t = \bigcap_{l=1}^{\ell} \bigcap_{\alpha \in \Gamma_l} E^l_\alpha. \tag{8.3} \]

More generally, let \(V\) be an \(n\)-dimensional combinatorial space of \(A^{<\mathbb{N}}\), and let \(\langle S_t : t \in V \rangle\) be a process in \((\Omega, \mathcal{F}, \mathbb{P})\). We say that \(\langle S_t : t \in V \rangle\) is \((\ell, r)\)-simplicial in \(V\) if the process \(\langle S_{t, V(t)} : t \in A^n \rangle\) is \((\ell, r)\)-simplicial.

**Remark 8.4.** In order to see the relevance of simplicial processes in this context note that if \(A, n, \ell\) and \(r\) are as in Definition 8.3 and \(\langle S_t : t \in A^n \rangle\) is an arbitrary \((\ell, r)\)-simplicial process, then there exist nonempty \(G \subseteq A^n\) and \(x \in A^n \setminus G\) such that the set \(G \cup \{x\}\) is \(\ell\)-separated and, moreover,

\[ \bigcap_{t \in G \cup \{x\}} S_t = \bigcap_{t \in G} S_t. \]

In particular, the events \(\langle S_t : t \in A^n \rangle\) cannot be independent.

### 8.2. The main result.

The following theorem—which is the main result in this section—complements Theorems 1.6 and 7.2 and completes the analysis of correlations of stationary processes. (We recall that the notion of stationarity for arbitrary correlations is given in Definition 5.7).

**Theorem 8.5.** Let \(k, \kappa, m\) be positive integers with \(k, \kappa \geq 2\) and \(\kappa \leq k^m\), and let \(\varepsilon, \sigma, \eta > 0\) such that

\[ \varepsilon \leq 1 - \frac{1}{2\kappa}, \quad \frac{\sigma}{2\kappa}, \quad \text{and } \eta \leq \frac{\sigma}{4^{\kappa-1}}. \tag{8.4} \]

Also let \(A\) be a set with \(|A| = k\), let \(n > m\) be an integer, and let \(\langle D_t : t \in A^n \rangle\) be an \(\eta\)-stationary process in a probability space \((\Omega, \mathcal{F}, \mathbb{P})\) such that \(|\mathbb{P}(D_t) - \varepsilon| \leq \eta\) for every \(t \in A^n\). Then, either
(i) for every nonempty $G \subseteq A^n$ with cardinality at most $\kappa$ and whose type $\tau(G)$ has dimension at most $m$ we have

\[
\left| \mathbb{P}\left( \bigcap_{t \in G} D_t \right) - \varepsilon^{[G]} \right| \leq \sigma,
\]

(ii) or $\langle D_t : t \in A^n \rangle$ correlates with a simplicial process when restricted on a large subspace; precisely, there exist $\ell \in [m]$ with the following property. If $\mathbf{r} = (r_1, \ldots, r_\ell)$ is an $\ell$-tuple of positive integers with $r := \sum_{i=1}^\ell r_i \leq n - m$, then there exist an $r$-dimensional combinatorial subspace $V$ of $A^n$ and a process $\langle S_t : t \in V \rangle$ in $(\Omega, \mathcal{F}, \mathbb{P})$ which is $(\ell, \mathbf{r})$-simplicial in $V$ such that for every $t \in V$ we have

\[
\mathbb{P}(S_t) \geq \frac{\varepsilon^{\kappa-1}}{4\kappa} \quad \text{and} \quad \mathbb{P}(D_t \mid S_t) \geq \varepsilon + \frac{\sigma}{4\kappa-1}.
\]

We have already pointed out that the proof of Theorem 8.5 is conceptually similar to the proofs of Theorems 1.6 and 7.2. More precisely, it relies on the following version of Propositions 3.7 and 7.3 which, in turn, is based on the higher-dimensional extensions of the archetypical “projection” $t^{\beta-\alpha}$. These extensions are presented in Subsection 8.3. (See Definition 5.10 for the notions of pseudorandomness, supercorrelation and subcorrelation which appear below.)

**Proposition 8.6.** Let $A$ be a finite set with $|A| \geq 2$, let $n, p, \ell$ be positive integers with $p + 1 \leq |A|^n$, let $0 < \eta, \varepsilon \leq 1$, and let $\langle D_t : t \in A^n \rangle$ be an $\eta$-stationary process in a probability space $(\Omega, \mathcal{F}, \mathbb{P})$ such that $\mathbb{P}(D_t) - \varepsilon \leq \eta$ for every $t \in A^n$. Let $\mathbf{t} = (t_1, \ldots, t_{p+1})$ be a tuple consisting of distinct elements of $A^n$ with $s(\mathbf{t}) = \ell$, set $G := \{t_1, \ldots, t_{p+1}\}$ and $H := \{t_1, \ldots, t_p\}$, and let $d$ denote the dimension of $\tau(G)$. Finally, let $0 < \theta, \sigma \leq 1$, let $\mathbf{r} = (r_1, \ldots, r_\ell)$ be an $\ell$-tuple of positive integers such that $r := \sum_{i=1}^\ell r_i \leq n - d$, and assume that the process $\langle D_t : t \in A^n \rangle$ is $(\tau(H), \theta)$-pseudorandom. Then there exist an $r$-dimensional combinatorial subspace $V$ of $A^n$ and a process $\langle S_t : t \in V \rangle$ in $(\Omega, \mathcal{F}, \mathbb{P})$ which is $(\ell, \mathbf{r})$-simplicial in $V$ with the following properties.

(i) For every $t \in V$ we have $|\mathbb{P}(S_t) - \varepsilon^n| \leq \theta$.

(ii) If $\langle D_t : t \in A^n \rangle$ is $(\tau(G), \sigma)$-supercorrelated, then for every $t \in V$ we have

\[
\mathbb{P}(D_t \mid S_t) \geq \varepsilon \left(1 + \frac{\sigma \varepsilon^{-1} - \theta}{\varepsilon \varepsilon^{-1} + \theta} \right).
\]

(iii) If $\langle D_t : t \in A^n \rangle$ is $(\tau(G), \sigma)$-subcorrelated, then for every $t \in V$ we have

\[
\mathbb{P}(D_t \mid S_t) \leq \varepsilon \left(1 - \frac{\sigma \varepsilon^{-1} - \theta}{\varepsilon \varepsilon^{-1} - \theta} \right).
\]

8.3. **Definitions/Notation.** This subsection is the analogue of Subsection 7.3. More precisely, let $A, n, p$ and $\ell$ be as in Proposition 8.6. Let

\[
\mathbf{t} = (t_1, \ldots, t_{p+1})
\]

\text{Note that, by Fact 6.3, if the type $\tau(G)$ of a nonempty set $G \subseteq A^n$ has dimension at most $m$, then $G$ is $\ell$-separated for some $\ell \leq m$.}
be an $\ell$-separated tuple consisting of distinct elements of $A^n$, and let $d$ be the dimension of $\tau := \tau(t)$. Also let $r = (r_1, \ldots, r_\ell)$ be a tuple of positive integers such that $r := \sum_{i=1}^\ell r_i \leq n - d$. We will define

- a set $J \subseteq [d]$ with $|J| = \ell$,
- $\beta_1, \ldots, \beta_\ell \in A$,
- an $r$-dimensional combinatorial subspace $V$ of $A^n$, and
- for every $j \in [p]$ a map $T_j : V \to A^n$.

These data are the combinatorial core of Theorem 8.5 and Proposition 8.6.

8.3.1. **Defining $J$ and $\beta_1, \ldots, \beta_\ell$.** We write the type $\tau(t) = (s_1, \ldots, s_{p+1})$ where $s_j = (s_j(1), \ldots, s_j(d)) \in A^d$ for every $j \in [p+1]$. Since $t$ is $\ell$-separated, by Fact 6.3 we have $s(\tau) = \ell$. Therefore, there exists $I \subseteq [d]$ with $|I| = \ell$ such that for every $j \in [p]$ there exists $i \in I$ satisfying $s_j(i) \neq s_{p+1}(i)$; let $J$ be the lexicographically least set with this property, write $J = \{i_1 < \cdots < i_\ell\}$, and set

\begin{equation}
\beta_1 := s_{p+1}(i_1), \ldots, \beta_\ell := s_{p+1}(i_\ell).
\end{equation}

Moreover, for every $j \in [p+1]$ and every $l \in [\ell+1]$ set

\begin{equation}
y_j^l = \begin{cases}
(s_j(i_{l-1}+1), \ldots, s_j(i_l)) & \text{if } l \in [\ell], \\
(s_j(i_{\ell}+1), \ldots, s_j(d)) & \text{if } l = \ell+1,
\end{cases}
\end{equation}

where $i_0 = 0$ and with the convention that $y_j^{\ell+1}$ is the empty sequence if $i_\ell = d$. Notice that $s_j = y_j^1 \cap \cdots \cap y_j^{\ell+1}$ for every $j \in [p+1]$.

8.3.2. **Defining $V$ and the maps ($T_j : j \in [p]$).** Set

\begin{equation}
V := \{y_{p+1}^{i_0} \cap z_1^{1} \cap \cdots \cap y_{p+1}^{i_\ell} \cap z_\ell^{\ell+1} : z_1 \in A^{r_1}, \ldots, z_\ell \in A^{r_\ell}\}
\end{equation}

and observe that $V$ is an $r$-dimensional combinatorial subspace of $A^n$. Finally, for every $j \in [p]$ we define $T_j : V \to A^n$ by the rule

\begin{equation}
T_j(y_{p+1}^1 \cap z_1^{1} \cap \cdots \cap y_{p+1}^{i_\ell} \cap z_\ell^{\ell+1}) = y_j^1 \cap \beta_1 \to s_j(i_1) \cdots \cap y_j^{i_\ell} \cap \beta_\ell \to s_j(i_\ell) \cdots \cap y_j^{\ell+1}
\end{equation}

with the convention $t^{\alpha - \alpha} = t$ for every $t \in A^{< N}$ and every $\alpha \in A$.

8.3.3. **Basic properties.** We close this subsection with the following analogue of Fact 7.4. The proof is straightforward.

**Fact 8.7.** Let $\ell, r, V$ and $\langle T_j : j \in [p]\rangle$ be as above.

(i) For every $t \in V$ and every $1 \leq i_1 < \cdots < i_q \leq p$ we have

\begin{equation}
\tau((T_{i_1}(t), \ldots, T_{i_q}(t), t)) = \tau((t_{i_1}, \ldots, t_{i_q}, t_{p+1}))
\end{equation}

and

\begin{equation}
\tau((T_{i_1}(t), \ldots, T_{i_q}(t))) = \tau((t_{i_1}, \ldots, t_{i_q})).
\end{equation}
(ii) Let $\langle D_t : t \in A^n \rangle$ be a stochastic process in a probability space $(\Omega, \Sigma, \mu)$. Let $\langle S_t : t \in V \rangle$ be a process of the form $S_t = \bigcap_{j=1}^p E_{T_j(t)}$ where for every $j \in [p]$ either $\langle E_{T_j(t)} : t \in V \rangle = \langle D_{T_j(t)} : t \in V \rangle$ or $\langle E_{T_j(t)} : t \in V \rangle = \langle D_{cT_j(t)} : t \in V \rangle$.

Then the process $\langle S_t : t \in V \rangle$ is $(\ell, r)$-simplicial in $V$.

8.4. Proof of Proposition 8.6. Let $V$ and $\langle T_j : j \in [p] \rangle$ be the data obtained in Subsection 8.3 for the $\ell$-separated tuple $t$ and the tuple $r$. We define $\langle S_t : t \in V \rangle$ by setting $S_t = \bigcap_{j=1}^p D_{T_j(t)}$ for every $t \in V$. By part (ii) of Fact 8.7, we see that the process $\langle S_t : t \in V \rangle$ is $(\ell, r)$-simplicial. Moreover, by (8.15) and our assumption that $\langle D_t : t \in A^n \rangle$ is $(\tau(H), \theta)$-pseudorandom, we have

$$|\mathbb{P}(S_t) - \varepsilon^p| \leq \theta$$

for every $t \in V$; that is, part (i) of the theorem holds true. The rest of the proof is identical to that of Proposition 7.3.

8.5. Proof of Theorem 8.5. Follows arguing precisely as in the proof of Theorem 7.2 using Proposition 8.6 and the material in Subsection 8.3.

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