The Space of Compatible Full Conditionals is a Unimodular Toric Variety

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Abstract

The set of all \( m \)-tuples of compatible full conditional distributions on discrete random variables is an algebraic set whose defining ideal is a unimodular toric ideal. We identify the defining polynomials of these ideals with closed walks on a bipartite graph. Our algebraic characterization provides a natural generalization of the requirement that compatible conditionals have identical odds ratios and holds regardless of the patterns of zeros in the conditional arrays.

Key words: Odds ratios, Toric ideal, Unimodular

1 Introduction

Statisticians have long been interested in combining marginal and conditional distributions in order to completely specify a joint distribution. Such models appear in spatial statistics (Besag, 1974), analysis of contingency tables (Bishop et al., 1975), Bayesian prior elicitation contexts (O’Hagan, 1998), expert systems (Cowell et al., 1999), statistical disclosure limitation (Slavkovic, 2004), and generally in any area of applied statistics where one wishes to build global statistical models from local information about subcollections of random variables.

Two fundamental theoretical questions which have been addressed in the literature are the compatibility of conditionals and marginals, and the uniqueness of the joint distribution when it exists. A collection of conditionals and marginals are compatible if there exists a joint distribution with these conditionals and marginals. The first major result along these lines is the Hammersley-Clifford theorem (Besag, 1974) which establishes a connection between the positive joint probability distribution and the full conditionals: it provides...
sufficient conditions for compatibility of distributions in the setting of Markov-random fields. In the discrete case, the Hammersley-Clifford theorem applies when there are no cells with zero probability. Questions about the uniqueness of the joint distribution given a compatible collection of marginals and conditionals were addressed in Gelman and Speed (1993) and the results presented there were subsequently generalized in Arnold et al. (1999). Arnold et al. (1999) also address the question of whether the given set of densities are compatible; they describe a variety of compatibility conditions for the case of finitely many discrete random variables and give algorithms for checking the compatibility of a set of conditional probability densities.

Let $X_1, \ldots, X_n$ be discrete random variables where the set of possible states for each $X_i$ is in the set of integers $[d_i] := \{1, 2, \ldots, d_i\}$. The joint probability distribution $p(x) = p(X_1 = i_1, \ldots, X_n = i_n)$ can be represented by a $d_1 \times \cdots \times d_n$ array of nonnegative real numbers that sum to one. We define a full conditional as a conditional multi-dimensional array $p(x_A | x_B)$ where $A \cup B = [n]$ and $A \cap B = \emptyset$ so that the conditional array depends on all the random variables. The precise statement of the compatibility problem for full conditionals is the following:

**Problem 1** Given partitions $(A_1, B_1), \ldots, (A_m, B_m)$ of $[n]$ and arrays $C^1 \ldots C^m$ each of format $d_1 \times \cdots \times d_n$, when is there a joint probability distribution $p(x)$ such that $p(x_A | x_B) = C^i$ for all $i$?

The compatibility problem makes sense for any set of conditionals and marginals, but we shall see a particularly simple combinatorial solution when we restrict attention to compatibility among full conditionals.

Given fixed values of the conditional and marginal arrays, the problem of determining the existence of a joint probability distribution can be phrased as a feasibility question for linear programs. Since marginal and conditional arrays are defined by linear and linear-fractional constraints in terms of the joint probabilities, the existence of a joint given the conditionals can be decided by determining whether or not a certain polytope is nonempty (Arnold et al., 1999). While this solution has many attractive features (the most important of which is the ability to include extra parameters to account for data that are “almost compatible” (Arnold et al., 1999)), it does not give an intrinsic characterization of those sets of arrays which correspond to compatible conditional distributions. In particular, using linear programming to determine whether or not a set of conditionals are compatible requires the introduction of extra parameters.

In this paper, we develop intrinsic methods for determining compatibility of conditional matrices, in the presence of zero cell entries. We describe explicit algebraic restrictions on the arrays $C^1, \ldots, C^m$ which guarantee that the pre-
scribed full conditional distributions are compatible. The results can be interpreted statistically in terms of conditions that replace generalized odds ratios to allow for arbitrary patterns of zeros. In this sense, our results provide a key link between the statistical notions of odds ratios and tools from algebraic geometry.

Our main result is the following theorem.

**Theorem 2** An \( m \)-tuple of \( d_1 \times \cdots \times d_n \) arrays \( C^1, \ldots, C^m \) are compatible full conditionals for the partitions \( (A_1, B_1), \ldots, (A_m, B_m) \) if and only if they satisfy the following four conditions:

1. \( C^i_{j_1 \ldots j_n} \geq 0 \) for all \( i \) and \( j_1, j_2, \ldots, j_n \).
2. If \( C^i_{j_1 \ldots j_n} = 0 \) for some \( i_0 \), then \( C^i_{j_1 \ldots j_n} = 0 \) for all \( i \).
3. For each \( i \) the \( B_i \) margin of \( C^i \) is an array of all ones.
4. The entries of the arrays \( C^i \) satisfy polynomial conditions which are in
   bijection with the induced circuits of a bipartite graph that depends on the
   \( d_1, d_2, \ldots, d_n \) and \( (A_1, B_1), \ldots, (A_m, B_m) \).

The outline for the rest of the paper is as follows. In the next section, we provide a detailed description of our results in the case of two random variables. We give precise defining relations for the space of compatible conditionals which generalize the condition that odds ratios of consistent conditionals should be the same. In particular, we give necessary and sufficient polynomial conditions for two conditionals to be compatible. In the third section, we describe the theoretical tools from combinatorial commutative algebra (Sturmfels, 1995) and discrete optimization (Schrijver, 1998) which we will need to prove Theorem 2. In the fourth section we provide a general proof that the space of compatible full conditionals is a unimodular toric variety. The fifth section is devoted to a few trivariate examples to illustrate the main theorem.

## 2 Two random variables

In this section, we give a detailed explanation of our algebraic results in the case of the bivariate compatible conditional problem. We refer the reader to (Cox et al., 1996) for the basics of algebraic geometry. We relegate technical definitions and proofs of the main theorems to the third and fourth sections. We denote the two potential conditional arrays by

\[
C_{ij} = p(X_1 = i | X_2 = j) = \frac{p(X_1 = i, X_2 = j)}{p(X_2 = j)} \quad \text{and}
\]
\[ D_{ij} = p(X_2 = j | X_1 = i) = \frac{p(X_1 = i, X_2 = j)}{p(X_1 = i)}. \]

This yields the following parametrization of the space of compatible conditionals.

**Proposition 3** Two nonnegative arrays \( C \) and \( D \) represent compatible conditionals for two discrete random variables if and only if there are nonnegative parameter arrays \( P_{ij}, u_i, v_j \) such that

\[ C_{ij} = P_{ij}v_j \quad \text{and} \quad D_{ij} = P_{ij}u_i, \]

\[ \sum_i C_{ij} = 1 \quad \text{for all} \ j \quad \text{and} \quad \sum_j D_{ij} = 1 \quad \text{for all} \ i. \]

Statistically, the array \( P_{ij} \) represent a joint distribution which has \( C \) and \( D \) as its corresponding conditionals. Arrays \( u_i \) and \( v_j \) are marginals which when combined with appropriate conditionals give \( P_{ij} \). This representation says that the set of all arrays \( C \) and \( D \) which are compatible conditions is determined by allowing the arrays \( P, u \) and \( v \) to range over all nonnegative values which make \( C \) and \( D \) into conditional arrays. The given parametrization is by polynomial functions in the parameters, and hence there exist polynomials in the \( C_{ij} \) and \( D_{ij} \) which vanish if and only if the pair \( C \) and \( D \) belong to the (closure of the) space of compatible conditionals. In the literature of discrete optimization and toric ideals, the parametrization defined in Proposition 3 is known as the Lawrence lifting of the Segre variety. The equations which vanish on this parametrization are well understood.

**Theorem 4** (e.g. Sturmfels, 1995, Ch. 14) The defining equations for the Lawrence lifting of the Segre variety are in bijection with circuits in the complete bipartite graph \( K_{d_1,d_2} \). Each circuit of length \( 2r \) defines a binomial of degree \( 2r \) in the \( C_{ij} \) and \( D_{ij} \).

This bijection is described as follows. Let \( (i_1, j_1, i_2, j_2, \ldots, i_r, j_r, i_1) \) be a circuit in \( K_{d_1,d_2} \) of length \( 2r \). This circuit produces a binomial of degree \( 2r \):

\[ C_{i_1j_1}D_{i_2j_2} \cdots D_{i_1j_r} = C_{i_1j_1}C_{i_2j_1}D_{i_2j_2} \cdots C_{i_1j_r}. \]

Note that the degree four relations which are produced by this construction are precisely the relations which say that all the odds ratios of the matrices \( C \) and \( D \) must be the same. For example, the circuit of length 4 in \( K_{2,2} \) produces the binomial:

\[ C_{11}C_{22}D_{12}D_{21} - C_{12}C_{21}D_{11}D_{22}. \]

For positive \( C \) and \( D \), this binomial is zero if and only if the following odds ratios are equivalent:

\[ \frac{C_{11}C_{22}}{C_{12}C_{21}} = \frac{D_{11}D_{22}}{D_{12}D_{21}}. \]
It is well known that for positive matrices $C$ and $D$, compatibility of conditional matrices is equivalent to equality of the odds ratios (Arnold et al., 1999). In algebro-geometric terms, this says that the degree four polynomials produced by the circuit construction define the toric variety of interest set theoretically in the strictly positive orthant. However, this statement becomes false when some of the entries in $C$ and $D$ are allowed to be zero. These polynomials, in addition to addressing the compatibility of arrays $C$ and $D$ in the presence of zero entries, also should be useful for determining the similarity of linear contrasts in the $C$ and $D$ matrices when there are zero cells (Bishop et al., 1975). Thus, these polynomials could prove useful for the analysis of incomplete contingency tables.

Example 5 Consider the matrices

$$C = \begin{pmatrix}
\frac{1}{2} & \frac{1}{2} & 0 \\
0 & \frac{1}{2} & \frac{1}{2} \\
\frac{1}{2} & 0 & \frac{1}{2}
\end{pmatrix} \quad \text{and} \quad D = \begin{pmatrix}
\frac{1}{3} & 0 & \frac{2}{3} \\
0 & \frac{1}{3} & \frac{2}{3} \\
\frac{1}{3} & 0 & \frac{2}{3}
\end{pmatrix}.$$  

All of the odds ratios of these matrices are the same: they are either simultaneously zero or undefined. In other words, they satisfy all 9 polynomials which are described by circuits in $K_{3,3}$ of length 4. However, the matrices do not represent compatible conditionals since the binomial given by the length 6 circuit $(1, 1, 2, 2, 3, 3, 1)$:

$$C_{11}C_{22}C_{33}D_{12}D_{23}D_{31} - C_{12}C_{23}C_{31}D_{11}D_{22}D_{33}$$

is equal to $1/108$ and not zero. These incompatible conditionals appear in Arnold et al. (1999).

These polynomials give a parameter free method for determining if two arbitrary conditional arrays are compatible.

To make the connection to the next section more lucid, we will introduce an alternate graph theoretic representation of the circuit binomials described above. This more complicated description is the one that generalizes for arbitrary sets of full conditionals in the next section.

Definition 6 We define the bipartite graph $G_{d_1,d_2}$ to be the graph with $d_1d_2 + d_1 + d_2$ vertices as follows.

1. There are $d_1d_2$ vertices labeled $v_{ij}$ where $i \in [d_1]$ and $j \in [d_2]$.
2. There are $d_1$ vertices labeled $w_k$ where $k \in [d_1]$.
3. There are $d_2$ vertices labeled $u_l$ where $l \in [d_2]$.
4. There is an edge between $v_{ij}$ and $w_k$ if and only if $i = k$.
5. There is an edge between $v_{ij}$ and $u_l$ if and only if $j = l$. 

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There are no other edges in $G_{d_1,d_2}$.

The graph $G_{2,3}$ is pictured in Figure 1. Associated to any circuit in $G_{d_1,d_2}$ of length $2r$ we get a binomial of degree $r$ as follows. After permuting indices, circuits in $G_{d_1,d_2}$ have the form

$$(i_1j_1, i_1, i_2j_2, j_2, i_2j_2, \ldots, j_r, i_rj_r)$$

where $i_1 = i_r$ and $j_1 = j_r$. From this circuit we recover the binomial

$$C_{i_1j_1}D_{i_2j_2} \cdots D_{i_1j_r} - D_{i_1j_1}C_{i_2j_2} \cdots C_{i_1j_r},$$

which was the same binomial from our previous description. Putting all these ideas together we have the following special case of Theorem 2.

**Corollary 7** A pair of matrices $C$ and $D$ are compatible conditionals if and only if they satisfy the following properties:

1. $C \geq 0$ and $D \geq 0$,
2. for all $i$ and $j$, $C_{ij} = 0$ if only if $D_{ij} = 0$,
3. $\sum_i C_{ij} = 1$ for all $j$ and $\sum_j D_{ij} = 1$ for all $i$, and
4. $C$ and $D$ satisfy the circuit polynomials associated to the circuits in the graph $G_{d_1,d_2}$.

## 3 Unimodular toric varieties

In this section, we will describe the mathematical objects that we will need to describe the space of compatible conditionals. First, we show that the space of compatible conditionals is a parametrized algebraic set. The parametrization is given by monomials and this leads naturally to the definition of toric
varieties and toric ideals. Then we introduce the notions of graphical unimodular matrices and unimodular toric varieties, which are the natural algebraic objects for representing the space of compatible conditionals.

A key preliminary observation is the following:

**Theorem 8** A set of full conditional matrices $C^1, \ldots, C^m$ are compatible conditionals corresponding to the partitions $(A_1, B_1), \ldots, (A_m, B_m)$ if and only if there are matrices of probabilities $U^1, \ldots, U^m$ of the appropriate sizes and a $d_1 \times \cdots \times d_n$ matrix of probabilities $P$ such that

$$P_{j_1 \ldots j_n} = C_{j_1 \ldots j_n}^{i} \cdot U_{j_k_1 \ldots j_{k_s}}^{i}$$

for all $i$ where $\{k_1 < \ldots < k_s\} = B_i$.

**PROOF.** This is the definition of compatibility together with Bayes’ rule. The matrices $U^i$ are then the “missing” $B_i$-marginal distributions. $\square$

Since we are interested in intrinsic characterizations of the compatible conditionals, it makes sense to rearrange this expression to deduce an equation for the $C^i$ in terms of the missing joint and conditional distributions $P$ and $U^i$. Thus we deduce the following corollary.

**Corollary 9** The set of all $m$-tuples of consistent full conditional arrays $C^1, \ldots, C^m$ is described parameterically as

$$C_{j_1 \ldots j_n}^{i} = P_{j_1 \ldots j_n}^{i} \cdot U_{j_k_1 \ldots j_{k_s}}^{i}.$$  

where the multiplication is componentwise.

At first glance, the space of conditionals we have described in Corollary 9 looks highly over-parameterized: shouldn’t the $U$ parameters be related to the $P$ parameters by marginalization?

We claim that we can first allow the parameters $P$ and $U^i$ to be arbitrary real arrays, consider the variety which is defined by this parameterization, and then intersect this variety with the appropriate product of probability simplices. The result of this procedure will be the same set of conditional arrays which is obtained by first restricting the parameters so that $U^i = p(x_B)$ and then considering the resulting parametrically described variety. To see this, observe that the only way that $C^i$ could be a conditional array is if $U^i$ is the corresponding marginal $p(x_B)$, which again follows by Bayes’ rule.

Let $p(x_{A_1}|x_{B_1})$ and $p(x_{A_2}|x_{B_2})$ be two full conditional matrices. Note that if $B_2 \subseteq B_1$, then they are compatible if and only if $p(x_{A_1}|x_{B_1})$ is obtained
from \( p(x_{A_2}|x_{B_2}) \) by conditioning on the variables \( A_2 \setminus A_1 \). Hence when determining the space of consistent conditionals for \( p(x_{A_1}|x_{B_1}), p(x_{A_2}|x_{B_2}), \ldots, p(x_{A_m}|x_{B_m}) \), we may assume that there are no containment relations among the \( A_i \) or \( B_i \). That is, the collection \( \Delta = \{B_1, \ldots, B_m\} \) can be taken to be the facets of a simplicial complex.

We will work in the polynomial ring
\[
R = \mathbb{R}[C_{j_1, \ldots, j_n}^i | 1 \leq i \leq m \text{ and } 1 \leq j_s \leq d_s, \text{ for all } s]
\]
which is a polynomial ring in \( D = md_1d_2\cdots d_n \) indeterminates. Each indeterminate in the ring \( R \) corresponds to an entry in one of the \( m \) conditional matrices. By Corollary 9, the algebraic variety in \( \mathbb{R}^D \), which, when intersected with a product of probability simplices, yields the space of consistent conditional, is given by the monomial parameterization
\[
C_{j_1, \ldots, j_n}^i = P_{j_1, \ldots, j_n}U_{j_{k_1}, \ldots, j_{k_s}}^i
\]
where \( \{k_1 < k_2 \cdots < k_s\} = B_i \) and the \( P \)'s and \( U \)'s are allowed to range over all of the real numbers.

**Definition 10** An algebraic variety which is given by a monomial parameterization is called a toric variety.

To describe the polynomials that define any toric variety amounts to understanding which products of the parameterizing monomials are equal to which other products of the parameterizing monomials. That is, the toric ideal which defines the toric variety is generated by binomials (differences of monomials).

**Definition 11** The toric ideal which is the vanishing ideal of the space of compatible conditionals is denoted \( I_\Delta \) which we call the compatibility ideal.

Furthermore, understanding which monomials are equal to each other can be studied by understanding the integer kernel of an associated matrix. A standard reference for the theory of toric ideals is [Sturmfels, 1995].

**Definition 12** Let \( A \in \mathbb{Z}^{d \times n} \) be an integer matrix. The toric variety associated to \( A \) is defined by the parameterization
\[
x_i = t_1^{a_{1i}}t_2^{a_{2i}}\cdots t_d^{a_{di}}.
\]
The toric ideal \( I_A \) which defines this parameterized variety is generated by the (infinitely many) binomials
\[
I_A = \langle x^u - x^v | Au = Av \rangle.
\]
The presentation of a toric ideal \( I_A \) in Definition 12 is given with infinitely
many ideal generators. However, the Hilbert basis theorem implies that only finitely many of these polynomials are actually needed in a generating set of the ideal. In many situations, the difficulty of studying toric ideals amounts to finding good combinatorial descriptions of this finite set of generators. In the case of the compatibility ideals $I_\Delta$ there is a simple combinatorial structure called unimodularity which can be used to describe the generators.

**Definition 13** A toric ideal $I_A$ is called unimodular if every reduced Gröbner basis of $I_A$ consists of squarefree binomials (i.e. there are no squared variables in the leading terms of each binomial in the reduced Gröbner basis). Equivalently, $I_A$ is unimodular if all the non-zero minors of $A$ have the same absolute value. Such a matrix $A$ is also called unimodular.

The following is an important special family of unimodular matrices (see, for example Schrijver (1998)).

**Proposition 14** Let $A$ be a matrix with the following properties:

(1) All the entries of $A$ are either zero or one,
(2) there are precisely two 1’s in each column of $A$, and
(3) there is a partition of the rows of $A$ into two sets $U$ and $V$ such that each column of $A$ has exactly one nonzero entry with row index $u \in U$ and one nonzero entry with row index $v \in V$.

Then $A$ is a unimodular matrix and is called a graphical unimodular matrix.

The reason for the name “graphical” is due to the fact that any graphical unimodular matrix $A$ is the vertex-edge incidence matrix of a bipartite graph $G_A$. The partition of the vertices of $G_A$ corresponds to the partition of the rows of $A$. Notice that vectors $v \in \ker(A)$ correspond to the union of cycles in the graph $G_A$.

**Definition 15** Let $G$ be any graph. An induced circuit of $G$ is a circuit of $G$ which does not have a chord (i.e. a “short cut”) in $G$.

In the complete bipartite graph $K_{3,3}$ all of the cycles of length 4 are induced, whereas none of the six cycles are induced because there is a chord cutting across them. Toric ideals which are presented by a graphical unimodular matrix have a simple combinatorial description for their minimal generators in terms of the induced circuits of the associated bipartite graph. A version of the following result can be found in (Aoki and Takemura, 2002).

**Proposition 16** Let $A$ be a graphical unimodular matrix, and $G_A$ the associated bipartite graph. Let the $U$ and $V$ be the partition of the vertices of $G_A$ where $U = \{u_1, \ldots, u_l\}$ and $V = \{v_1, \ldots, v_l\}$. For each circuit $c =$
(u_{i1}, v_{j1}, u_{i2}, v_{j2}, \ldots, u_{ir}, v_{jr}) we associate the circuit binomial
\[ f_c = x_{u_{i1}v_{j1}}x_{u_{i2}v_{j2}} \cdots x_{u_{ir}v_{jr}} - x_{u_{i2}v_{j1}}x_{u_{i3}v_{j2}} \cdots x_{u_{ir}v_{jr}}. \]

Then \( I_A \) is minimally generated by the circuit binomials corresponding to the induced circuits of \( G_A \). That is,
\[ I_A = \langle f_c \mid c \text{ is an induced circuit of } G_A \rangle. \]

4 The main algebraic result

Now that we have reviewed all the mathematical facts we need about unimodular toric ideals, we are in a position to prove that the compatibility ideal is a unimodular toric ideal. To do this, we construct the matrix that represents this toric ideal.

Let \( d = (d_1, \ldots, d_n) \) be the integer vector corresponding to the dimensions of the tables let \( \Delta = \{ B_1, \ldots, B_m \} \) be the simplicial complex which represents the compatible conditionals. For each facet \( B_i \) of \( \Delta \), denote by \( d_{B_i} \) the vector \( d_{B_i} = (d_{k1}, \ldots, d_{ks}) \) where \( B_i = \{ k_1, \ldots, k_s \} \) and denote by \( D_{B_i} \) the product
\[ D_{B_i} = \prod_{k_i \in F_i} d_{k_i}. \]

We denote by \( A_{\Delta,d} \) the matrix that represents the compatibility ideal corresponding to this data. This matrix is a
\[ (d_1d_2 \cdots d_n + \sum_{i=1}^m d_{B_i}) \times md_1d_2 \cdots d_n \]
matrix. The rows are naturally grouped into \( m+1 \) blocks and the columns are naturally grouped into \( m \) blocks. Each column is labelled by an integer \( i \in [m] \) and an \( n \)-tuple \( (j_1, \ldots, j_n) \in [d_1] \times [d_2] \times \cdots \times [d_n] \). The first block of rows is labelled by the integer 0 and \( n \)-tuple \( (j_1, \ldots, j_n) \in [d_1] \times [d_2] \times \cdots \times [d_n] \). Each of the remaining blocks of rows is labelled by an integer \( i \) and an \( |B_i| \)-tuple \( (j_{k1}, \ldots, j_{ks}) \in [d_{k1}] \times \cdots \times [d_{ks}] \) where \( B_i = \{ k_1, \ldots, k_s \} \).

The entries of this matrix are all ones and zeros. There is a one in a particular entry with row indexed by the data \( i^1, (j_{1}^1, \ldots, j_{d}^1) \) and column indexed by the data \( i^2, (j_{1}^2, \ldots, j_{d}^2) \) if and only if it satisfies the following rules:

(1) if the row label is 0 and \( (j_{1}^1, \ldots, j_{d}^1) = (j_{1}^2, \ldots, j_{d}^2) \) or
(2) if the row label is $i^1 > 0$, $i^1 = i^2$, and $(j_{k_1}^1, \ldots, j_{k_s}^1) = (j_{k_1}^2, \ldots, j_{k_s}^2)$ where $B_i = (k_1, \ldots, k_s)$.

**Theorem 17** The matrix $A_{\Delta,d}$ is a graphical unimodular matrix and it represents the toric variety of the space of compatible conditionals. Hence, the compatibility ideal is generated by the induced circuit binomials in the associated bipartite graph.

**PROOF.** To see that $A_{\Delta,d}$ represents the toric variety of compatible conditionals amounts to identifying the rows of the matrix with a parameter, and the columns of the matrix with an entry in the conditional matrix. The labelling for the columns $i, (j_1, \ldots, j_n)$ naturally corresponds to the indeterminate $C_{i^1 \ldots j_n}$. A row in the first block, with $i = 0$, corresponds to the $P_{j_1, \ldots, j_n}$ parameters. A row in any of the blocks with $i > 0$ corresponds to the parameter $U_{i^1 \ldots j_n}^i$. Note that be the description of the matrix $A_{\Delta,d}$, in the column corresponding to $C_{i^1 \ldots j_n}^i$ there are precisely two nonzero entries: one in the first block which corresponds to $P_{j_1, \ldots, j_n}$ and one in block $i$, which corresponds to $U_{i^1 \ldots j_n}^i$. Thus, this represents the parameterization

$$C_{i^1 \ldots j_n}^i = P_{j_1 \ldots j_n} U_{i_k^1 \ldots j_k^s}^i$$

as desired.

Now we wish to show the unimodularity of $A_{\Delta,d}$, however, this is an immediate consequence of the preceding argument. Each column has precisely two ones: one with index $i = 0$ and one with index $i > 0$. Hence, the value of $i$ determines the partition of the rows to deduce the structure of the underlying bipartite graph. \(\square\)

Now we will explicitly describe the graph encoded by the matrix $A_{\Delta,d}$.

**Definition 18** The graph $G_{\Delta,d}$ associated to the compatibility ideal has

$$(d_1 d_2 \cdots d_n + \sum_{i=1}^m d_{F_i})$$

vertices and $md_1 d_2 \cdots d_n$ edges: these are the number of rows and columns of $A_{\Delta,d}$ respectively. The vertices are partitioned into two class: those labelled with $i = 0$ and some $(j_1, \ldots, j_n)$, and those labelled with $i > 0$ and some $(j_{k_1}, \ldots, j_{k_s})$. Vertices with $i = 0$ are incident only to vertices with $i > 0$ and conversely. In particular, the vertex $i = 0, (j_1^1, \ldots, j_n^1)$ is incident to $i > 0$, $(j_{k_1}^2, \ldots, j_{k_s}^2)$ if and only if $(j_{k_1}^1, \ldots, j_{k_s}^1) = (j_{k_1}^2, \ldots, j_{k_s}^2)$.

To close this section, we now provide the proof of Theorem 2.
PROOF of Theorem 2: Conditions (1)-(3) of Theorem 2 are clearly necessary. Condition (4) is simply the expression of Theorem 17 above. What remains to show is that this condition is, in fact, sufficient. Suppose that a particular realization of conditional arrays $C^1, \ldots, C^m$ satisfy conditions (1)-(4). We wish to show that there exists parameter matrices $P$ and $U^1, \ldots, U^m$, which have $C^1, \ldots, C^m$ as their image. By conditions (1), (3), and (4) together with a result of Geiger, et al. (2002), this can happen if and only if the columns of $A_{\Delta,d}$ which are indexed by the support of $C^1, \ldots, C^m$ form a nice facial subset of the columns of $A_{\Delta,d}$. The nice facial subsets of the columns of $A_{\Delta,d}$ are precisely those sets of columns which are obtained by taking a collection of rows $R$ of $A_{\Delta,d}$, and taking exactly those columns of $A_{\Delta,d}$ which have zeros in the rows $R$. However, the support of $C^1, \ldots, C^m$ is a nice facial subset of the columns of $A_{\Delta,d}$ because of condition (2): the rows of $A_{\Delta,d}$ to choose are precisely those rows labelled by $(0, (j_1, \ldots, j_n))$ for all the indices $(j_1, \ldots, j_n)$ where the $C^i_{j_1,\ldots,j_n}$ are collectively zero. □

5 Trivariate examples

We will conclude our paper with some examples on three binary random variables to illustrate the construction and the types of polynomials that appear.

Example 19 Consider the compatibility ideal associated with three binary random variables and suppose that $\Delta = \{\{1\}, \{2\}, \{3\}\}$. In other words, we are considering the compatibility of $p(x_1, x_2 | x_3), p(x_1, x_3 | x_2)$ and $p(x_2, x_3 | x_1)$. The matrix, $A_{\Delta,(2,2,2)}$ has 14 rows and 24 columns. It is the zero/one matrix

$$
\begin{pmatrix}
1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\
0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 
\end{pmatrix}
$$

The graph corresponding to this compatibility problem has 14 vertices: 8 of one type and 6 of the other, and 24 edges between them. It is pictured in Figure 2.

There are three types of induced circuits in this graph, yielding three types of binomial generators of the compatibility ideal $I_{\Delta}$. We use the indeterminates $C_{ijk}, D_{ijk}$ and $E_{ijk}$ to denote the corresponding entries in the conditional ar-
rays. For instance, the circuit \((111, c_1, 211, d_1, 111)\) in the graph is induced and from it we deduce that the quadratic binomial

\[ C_{111}D_{211} - D_{111}C_{211} \]

is a minimal generator of the compatibility ideal \(I_\Delta\). The induced circuit \((111, c_1, 221, e_2, 212, d_1, 111)\) produces the binomial

\[ C_{111}E_{221}D_{212} - C_{221}E_{212}D_{111} \]

which is also a minimal generator of \(I_\Delta\). In total, there are 12 induced circuits of length 4, 8 induced circuits of length 6, and 60 induced circuits of length 8. Hence \(I_\Delta\) is minimally generated by 12 quadrics, 8 cubics, and 60 quartics.

**Example 20** Consider the compatibility associated to 3 binary random variables with the simplicial complex \(\Delta = \{\{1, 2\}, \{1, 3\}, \{2, 3\}\}\). That is, we are considering the compatibility of \(p(x_1|x_2, x_3), p(x_1|x_3, x_2), p(x_2|x_3, x_1)\). Use the indeterminates \(C_{ijk}, D_{ijk}, E_{ijk}\) to represent entries in the corresponding conditional arrays. The ideal \(I_\Delta\) is generated by 28 binomials which fall into 4 equivalence classes modulo the symmetry of the cube. In this case, the binomial minimal generators of \(I_\Delta\) are in bijection with the circuit in the edge graph of the cube. The four different types of circuit on the cube are indicated in figure 3.

**Fig. 3. The combinatorial types of circuits on the cube**
These four types of circuits yield four different symmetry classes of binomials in the ideal $I_\Delta$ which are represented by the following binomials:

$$C_{111}C_{221}D_{121}D_{211} - C_{121}C_{211}D_{111}D_{221}$$
$$C_{111}D_{211}E_{221}D_{222}C_{212}E_{112} - C_{211}D_{221}E_{222}D_{212}C_{112}E_{111},$$
$$C_{111}D_{211}E_{222}C_{222}D_{122}E_{112} - C_{211}D_{221}E_{222}C_{122}D_{112}E_{111},$$
$$C_{111}D_{211}C_{221}E_{121}C_{122}D_{222}C_{212}E_{112} - C_{211}D_{221}C_{121}E_{122}C_{222}D_{212}C_{112}E_{111}.$$

Note that the first polynomial in this list has a natural statistical interpretation: it corresponds to the equality of the odds ratios for $p(x|y, z)$ and $p(y|x, z)$.

While there are statistical interpretations for some of the circuit polynomials which arise, understanding these polynomial expressions in relation to known statistical quantities remains an open problem. In particular, determining how tools from the analysis of contingency tables such as generalized odds ratios and linear contrast relate to the polynomial constraints we have derived could prove useful for inference on incomplete contingency tables.

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