Sample-Efficient Reinforcement Learning of Undercomplete POMDPs

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Abstract

Partial observability is a common challenge in many reinforcement learning applications, which requires an agent to maintain memory, infer latent states, and integrate this past information into exploration. This challenge leads to a number of computational and statistical hardness results for learning general Partially Observable Markov Decision Processes (POMDPs). This work shows that these hardness barriers do not preclude efficient reinforcement learning for rich and interesting subclasses of POMDPs. In particular, we present a sample-efficient algorithm, OOM-UCB, for episodic finite undercomplete POMDPs, where the number of observations is larger than the number of latent states and where exploration is essential for learning, thus distinguishing our results from prior works. OOM-UCB achieves an optimal sample complexity of $O(1/\varepsilon^2)$ for finding an $\varepsilon$-optimal policy, along with being polynomial in all other relevant quantities. As an interesting special case, we also provide a computationally and statistically efficient algorithm for POMDPs with deterministic state transitions.

1 Introduction

In many sequential decision making settings, the agent lacks complete information about the underlying state of the system, a phenomenon known as partial observability. Partial observability significantly complicates the tasks of reinforcement learning and planning, because the non-Markovian nature of the observations forces the agent to maintain memory and reason about beliefs of the system state, all while exploring to collect information about the environment. For example, a robot may not be able to perceive all objects in the environment due to occlusions, and it must reason about how these objects may move to avoid collisions [Cassandra et al., 1996]. Similar reasoning problems arise in imperfect information games [Brown and Sandholm, 2018], medical diagnosis [Hauskrecht and Fraser, 2000], and elsewhere [Rafferty et al., 2011]. Furthermore, from a theoretical perspective, well-known complexity-theoretic results show that learning and planning in partially observable environments is statistically and computationally intractable in general [Papadimitriou and Tsitsiklis, 1987, Mundhenk et al., 2000, Vlassis et al., 2012, Mossel and Roch, 2005].

The standard formulation for reinforcement learning with partial observability is the Partially Observable Markov Decision Process (POMDP), in which an agent operating on noisy observations makes decisions that influence the evolution of a latent state. The complexity barriers apply for this model, but they are of a worst case nature, and they do not preclude efficient algorithms for interesting sub-classes of POMDPs. Thus we ask:

Can we develop efficient algorithms for reinforcement learning in large classes of POMDPs?
This question has been studied in recent works [Azizzadenes heli et al., 2016, Guo et al., 2016], which incorporate a decision making component into a long line of work on “spectral methods” for estimation in latent variable models [Hsu et al., 2012, Song et al., 2010, Anandkumar et al., 2012, 2014], including the Hidden Markov Model. Briefly, these estimation results are based on the method of moments, showing that under certain assumptions the model parameters can be computed by a decomposition of a low-degree moment tensor. The works of Azizzadenesheli et al. [Azizzadenesheli et al., 2016] and Guo et al. [Guo et al., 2016] use tensor decompositions in the POMDP setting and obtain sample efficiency guarantees. Neither result considers a setting where strategic exploration is essential for information acquisition, and they do not address one of the central challenges in more general reinforcement learning problems.

Our contributions. In this work, we provide new sample-efficient algorithms for reinforcement learning in finite POMDPs in the undercomplete regime, where the number of observations is larger than the number of latent states. This assumption is quite standard in the literature on estimation in latent variable models [Anandkumar et al., 2014]. Our main algorithm OOM-UCB uses the principle of optimism for exploration and uses the information gathered to estimate the Observable Operators induced by the environment. Our main result proves that OOM-UCB finds a near optimal policy for the POMDP using a number of samples that scales polynomially with all relevant parameters and additionally with the minimum singular value of the emission matrix. Notably, OOM-UCB finds an $\varepsilon$-optimal policy at the optimal rate of $O(1/\varepsilon^2)$.

While OOM-UCB is statistically efficient for this subclass of POMDPs, we should not expect it to be computationally efficient in general, as this would violate computational barriers for POMDPs. However, in our second contribution, we consider a further restricted subclass of POMDPs in which the latent dynamics are deterministic and where we provide both a computationally and statistically efficient algorithm. Notably, deterministic dynamics are still an interesting subclass due to that, while it avoids computational barriers, it still does not mitigate the need for strategic exploration. We prove that our second algorithm has sample complexity scaling with all the relevant parameters as well as the minimum $\ell_2$ distance between emission distributions. This latter quantity replaces the minimum singular value in the guarantee for OOM-UCB and is a more favorable dependency.

We provide further motivation for our assumptions with two lower bounds: the first shows that the over-complete setting is statistically intractable without additional assumptions, while the second necessitates the dependence on the minimum singular value of the emission matrix. In particular, under our assumptions, the agent must engage in strategic exploration for sample-efficiency. As such, the main conceptual advance in our line of inquiry over prior works is that our algorithms address exploration and partial observability in a provably efficient manner.

1.1 Related work

A number of computational barriers for POMDPs are known. If the parameters are known, it is PSPACE-complete to compute the optimal policy, and, furthermore, it is NP-hard to compute the optimal memoryless policy [Papadimitriou and Tsitsiklis, 1987, Vlassis et al., 2012]. With regards to learning, Mossel and Roch [Mossel and Roch, 2005] provided an average case computationally complexity result, showing that parameter estimation for a subclass of Hidden Markov Models (HMMs) is at least as hard as learning parity with noise. This directly implies the same hardness result for parameter estimation in POMDP models, due to that an HMM is just a POMDP with a fixed action sequence. On the other hand, for reinforcement learning in POMDPs (in particular, finding a near optimal policy), one may not need to estimate the model, so this lower bound need not directly imply that the RL problem is computational intractable. In this work, we do provide a lower bound showing that reinforcement learning in POMDPs is both statistically and computationally intractable (Propositions 1 and 2).

On the positive side, there is a long history of work on learning POMDPs. [Even-Dar et al., 2005] studied
POMDPs without resets, where the proposed algorithm has sample complexity scaling exponentially with a certain horizon time, which is not possible to relax without further restrictions. \cite{Ross2008, Poupart2008} proposed to learn POMDPs using Bayesian methods; PAC or regret bounds are not known for these approaches.

Closest to our work are POMDP algorithms based on spectral methods \cite{Guo2016, Azizzadenesheli2016}, which were originally developed for learning latent variable models \cite{Hsu2012, Anandkumar2012, Song2010, Sharan2017}. These works give PAC and regret bounds (respectively) for tractable subclasses of POMDPs, but, in contrast with our work, they make additional assumptions to mitigate the exploration challenge. In \cite{Guo2016}, it is assumed that all latent states can be reached with nontrivial probability with a constant number of random actions. This allows for estimating the entire model without sophisticated exploration. \cite{Azizzadenesheli2016} consider a special class of memoryless policies in a setting where all of these policies visit every state and take every action with non-trivial probability. As with \cite{Guo2016}, this restriction guarantees that the entire model can be estimated regardless of the policy executed, so sophisticated exploration is not required. We also mention that \cite{Guo2016, Azizzadenesheli2016} assume that both the transition and observation matrices are full rank, which is stronger than our assumptions. We do not make any assumptions on the transition matrix.

Finally, the idea of representing the probability of a sequence as products of operators dates back to multiplicity automata \cite{Schützenberger1961, Carlyle1971} and reappeared in the Observable Operator Model (OOMs) \cite{Jaeger2000} and Predictive State Representations (PSRs) \cite{Littman2002}. While spectral methods have been applied to PSRs \cite{Boots2011}, we are not aware of results with provable guarantees using this approach. It is also worth mentioning that any POMDP can be modeled as an Input-Output OOM \cite{Jaeger1998}.

## 2 Preliminaries

In this section, we define the partially observable Markov decision process, the observable operator model \cite{Jaeger2000}, and discuss their relationship.

**Notation.** For any natural number $n \in \mathbb{N}$, we use $[n]$ to denote the set $\{1, 2, \ldots, n\}$. We use bold upper-case letters $\mathbf{B}$ to denote matrices and bold lower-case letters $\mathbf{b}$ to denote vectors. $B_{ij}$ means the $(i, j)^{th}$ entry of matrix $\mathbf{B}$ and $(\mathbf{B})_i$ represents its $i^{th}$ column. For vectors we use $\| \cdot \|_p$ to denote the $\ell_p$-norm, and for matrices we use $\| \cdot \|_p$, $\| \cdot \|_1$ and $\| \cdot \|_F$ to denote the spectral norm, entrywise $\ell_1$-norm and Frobenius norm respectively. We denote by $\| \mathbf{B} \|_{p \rightarrow q} = \max_{\| \mathbf{v} \|_p \leq 1} \| \mathbf{B} \mathbf{v} \|_q$ the $p$-to-$q$ norm of $\mathbf{B}$. For any matrix $\mathbf{B} \in \mathbb{R}^{m \times n}$, we use $\sigma_{\min}(\mathbf{B})$ to denote its smallest singular value, and $\mathbf{B}^\dagger \in \mathbb{R}^{n \times m}$ to denote its Moore-Penrose inverse. For vector $\mathbf{v} \in \mathbb{R}^n$, we denote $\text{diag}(\mathbf{v}) \in \mathbb{R}^{n \times n}$ as a diagonal matrix where $[\text{diag}(\mathbf{v})]_{ii} = v_i$ for all $i \in [n]$. Finally, we use standard big-O and big-Omega notation $\tilde{O}(\cdot)$, $\Omega(\cdot)$ to hide only absolute constants which do not depend on any problem parameter, and notation $\tilde{O}(\cdot)$, $\Omega(\cdot)$ to hide only absolute constants and logarithmic factors.

### 2.1 Partially observable Markov decision processes

We consider an episodic tabular Partially Observable Markov Decision Process (POMDP), which can by specified as $\text{POMDP}(H, \mathcal{S}, \mathcal{A}, \mathcal{O}, \mathcal{T}, \mathcal{D}, r, \mu_1)$. Here $H$ is the number of steps in each episode, $\mathcal{S}$ is the set of states with $|\mathcal{S}| = S$, $\mathcal{A}$ is the set of actions with $|\mathcal{A}| = A$, $\mathcal{O}$ is the set of observations with $|\mathcal{O}| = O$, $\mathcal{T} = \{T_h\}_{h=1}^H$ specify the transition dynamics such that $T_h(\cdot | s, a)$ is the distribution over states if action $a$ is taken from state $s$ at step $h \in [H]$, $\mathcal{D} = \{D_h\}_{h=1}^H$ are emissions such that $D_h(\cdot | s)$ is the distribution over observations for state
s at step $h \in [H]$, $r = \{r_h : \mathcal{O} \rightarrow [0, 1]\}_{h=1}^H$ are the known deterministic reward functions$^1$ and $\mu_1(\cdot)$ is the initial distribution over states. Note that we consider nonstationary dynamics, observations, and rewards.

In a POMDP, states are hidden and unobserved to the learning agent. Instead, the agent is only able to see the observations and its own actions. At the beginning of each episode, an initial hidden state $s_1$ is sampled from initial distribution $\mu_1$. At each step $h \in [H]$, the agent first observes $o_h \in \mathcal{O}$ which is generated from the hidden state $s_h \in \mathcal{S}$ according to $\mathbb{P}(o_h|s_h)$, and receives the reward $r_h(o_h)$, which can be computed from the observation $o_h$. Then, the agent picks an action $a_h \in \mathcal{A}$, which causes the environment to transition to hidden state $s_{h+1}$, that is drawn from the distribution $\mathbb{P}(s_{h+1}|s_h, a_h)$. The episode ends when $o_H$ is observed.

A policy $\pi$ is a collection of $H$ functions $\{\pi_h : \mathcal{A} \rightarrow \mathcal{S}\}_{h=1}^H$, where $\mathcal{S} = (\mathcal{O} \times \mathcal{A})^{h-1} \times \mathcal{O}$ is the set of all possible histories of length $h$. We use $V^\pi \in \mathbb{R}$ to denote the value of policy $\pi$, so that $V^\pi$ gives the expected cumulative reward received under policy $\pi$:

$$V^\pi := \mathbb{E}_\pi \left[ \sum_{h=1}^H r_h(o_h) \right].$$

Since the state, action, observation spaces, and the horizon, are all finite, there always exists an optimal policy $\pi^*$ which gives the optimal value $V^* = \sup_{\pi} V^\pi$. We remark that, in general, the optimal policy of a POMDP will select actions based the entire history, rather than just the recent observations and actions. This is one of the major differences between POMDPs and standard Markov Decision Processes (MDPs), where the optimal policies are functions of the most recently observed state. This difference makes POMDPs significantly more challenging to solve.

**The POMDP learning objective.** Our objective in this paper is to learn an $\epsilon$-optimal policy $\tilde{\pi}$ in the sense that $V^\tilde{\pi} \geq V^* - \epsilon$, using a polynomial number of samples.

### 2.2 The observable operator model

We have described the POMDP model via the transition and observation distributions $T$, $\mathbb{P}$ and the initial distribution $\mu_1$. While this parametrization is natural for describing the dynamics of the system, POMDPs can also be fully specified via a different set of parameters: a set of operators $\{B_h(a, o) \in \mathbb{R}^{O \times S}\}_{h=1}^H$, action $a \in \mathcal{A}$ and observation $o \in \mathcal{O}$, and a vector $b_0 \in \mathbb{R}^O$.

In the undercomplete setting where $S \leq O$ and where observation probability matrices $\{\mathbb{P}_h \in \mathbb{R}^{O \times S}\}_{h=1}^H$ are all full column-rank, the operators $\{B_h(a, o)\}_{h,a,o}$ and vector $b_0$ can be expressed in terms of $(T, \mathbb{P}, \mu_1)$ as follows:

$$B_h(a, o) = \mathbb{P}_{h+1} T_h(a) \text{diag}(\mathbb{P}_h(o|\cdot)) \mathbb{P}_h^{-1}, \quad b_0 = \mathbb{P}_{1}\mu_1. \quad (1)$$

where we use the matrix and vector notation for $\mathbb{P}_h \in \mathbb{R}^{O \times S}$ and $\mu_1 \in \mathbb{R}^S$ here, such that $[\mathbb{P}_h]_{a,o} = \mathbb{P}_h(o|s)$ and $[\mu_1]_s = \mu_1(s)$. $T_h(a) \in \mathbb{R}^{S \times S}$ denotes the transition matrix given action $a \in \mathcal{A}$ where $[T_h(a)]_{s',s} = T_h(s'|s,a)$, and $\mathbb{P}_h(o|\cdot) \in \mathbb{R}^S$ denotes the $o$-th row in matrix $\mathbb{P}_h$ with $[\mathbb{P}_h(o|\cdot)]_s = \mathbb{P}_h(o|s)$. Using these matrices $B_h$, it can be shown that (Fact 18 in the appendix), for any sequence of $(o_H, \ldots, a_1, o_1) \in \mathcal{O} \times (\mathcal{A} \times \mathcal{O})^{H-1}$, we have:

$$\mathbb{P}(o_H, \ldots, o_1|a_{H-1}, \ldots, a_1) = e_{o_H}^T \cdot B_{H-1}(a_{H-1}, o_{H-1}) \cdots B_1(a_1, o_1) \cdot b_0. \quad (2)$$

Describing these conditional probabilities for every sequence is sufficient to fully specify the entire dynamical system. Therefore, as an alternative to directly learning $T$, $\mathbb{P}$ and $\mu_1$, it is also sufficient to learn operators

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$^1$Since rewards are observable in most applications, it is natural to assume the reward is a known function of the observation. While we study deterministic reward functions for notational simplicity, our results generalize to randomized reward functions. Also, we assume the reward is in $[0, 1]$ without loss of generality.
\{B_h(a, o)\}_{h,a,o} and vector \(b_0\) in order to learn the optimal policy. The latter approach enjoys the advantage that (2) does not explicitly involve latent variables. It refers only to observable quantities—actions and observations.

We remark that the operator model introduced in this section (which is parameterized by \{B_h(a, o)\}_{h,a,o} and \(b_0\)) bears significant similarity to Jaeger’s Input-Output Observable Operator Model (IO-OOM) [Jaeger, 2000], except a few minor technical differences.\(^2\) With some abuse of terminology, we also refer to our model as Observable Operator Model (OOM) in this paper. It is worth noting that Jaeger’s IO-OOMs are strictly more general than POMDPs [Jaeger, 2000] and also includes overcomplete POMDPs via relation different from (1).

Since our focus is on undercomplete POMDPs, we refer the reader to Jaeger [2000] for more details.

3 Main Results

We first state our main assumptions, which we motivate with corresponding hardness results in their absence. We then present our main algorithm, \(OOM-UCB\), along with its sample efficiency guarantee.

3.1 Assumptions

In this paper, we make following assumptions.

**Assumption 1.** We assume the POMDP is undercomplete, i.e. \(S \leq O\). We also assume the minimum singular value of the observation probability matrices \(\sigma_{\text{min}}(\mathcal{O}_h) \geq \alpha > 0\) for all \(h \in [H]\).

Both assumptions are standard in the literature on learning Hidden Markov Models (HMMs)—an uncontrolled version of POMDP [see e.g., Anandkumar et al., 2012]. The second assumption that \(\sigma_{\text{min}}(\mathcal{O}_h)\) is lower-bounded is a robust version of the assumption that \(\mathcal{O}_h \in \mathbb{R}^{O \times S}\) is full column-rank, which is equivalent to \(\sigma_{\text{min}}(\mathcal{O}_h) > 0\). Together, these assumption ensure that the observations will contain a reasonable amount of information about the latent states.

We do not assume that the initial distribution \(\mu_1\) has full support, nor do we assume the transition probability matrices \(T_h\) are full rank. In fact, Assumption 1 is not sufficient for identification of the system, i.e. recovering parameters \(T, \mathcal{O}, \mu_1\) in total-variance distance. Exploration is crucial to find a near-optimal policy in our setting.

We motivate both assumptions above by showing that, with absence of either one, learning a POMDP is statistically intractable. That is, it would require an exponential number of samples for any algorithm to learn a near-optimal policy with constant probability.

**Proposition 1.** For any algorithm \(\mathfrak{A}\), there exists an overcomplete POMDP \((S > O)\) with \(S\) and \(O\) being small constants, which satisfies \(\sigma_{\text{min}}(\mathcal{O}_h) = 1\) for all \(h \in [H]\), such that algorithm \(\mathfrak{A}\) requires at least \(\Omega(H^{-1})\) samples to ensure learning a \((1/4)\)-optimal policy with probability at least 1/2.

**Proposition 2.** For any algorithm \(\mathfrak{A}\), there exists an undercomplete POMDP \((S \leq O)\) with \(S\) and \(O\) being small constants, such that algorithm \(\mathfrak{A}\) requires at least \(\Omega(H^{-1})\) samples to ensure learning a \((1/4)\)-optimal policy with probability at least 1/2.

Proposition 1 and 2 are both proved by constructing hard instances, which are modifications of classical combinatorial locks for MDPs [Krishnamurthy et al., 2016]. We refer readers to Appendix B for more details.

3.2 Algorithms

We are now ready to describe our algorithm. Assumption 1 enables the representation of the POMDP using OOM with relation specified as in Equation (1). Our algorithm, Observable Operator Model with Upper
Algorithm 1 Observable Operator Model with Upper Confidence Bound (OOM-UCB)

1: **Initialize:** set all entries in a vector of counts \( n \in \mathbb{N}^O \), and in matrices of counts \( N_h(a, \bar{a}) \in \mathbb{N}^{O \times O} \), \( M_h(o, a, \bar{a}) \in \mathbb{N}^{O \times O} \) to be zero for all \((o, a, \bar{a}) \in \mathcal{O} \times \mathcal{A}^2\).

2: set confidence set \( \Theta_1 = \bigcap_{h \in [H]} \{ \theta \mid \sigma_{\min}(\hat{\Theta}_h) \geq \alpha \} \).

3: for \( k = 1, 2, \ldots, K \) do

4: compute the optimistic policy \( \pi_k \leftarrow \arg \max_a \max_{\theta \in \Theta_k} V^\pi(\hat{\theta}) \).

5: observe \( o_1 \), and set \( n \leftarrow n + e_{o_1} \).

6: \( b \leftarrow \bigcap_{h \in [H]} \{ \theta \mid \sigma_{\min}(\hat{\Theta}_h) \geq \alpha \} \cap \{ \theta \mid \| k \cdot b_0(\hat{\theta}) - n \|_2 \leq \beta_k \} \).

7: for \((h, a, \bar{a}) \in [H-1] \times \mathcal{A}^2 \) do

8: execute policy \( \pi_k \) from step 1 to step \( h - 2 \).

9: take action \( \bar{a} \) at step \( h - 1 \), and action \( a \) at step \( h \) respectively.

10: observe \((o_{h-1}, o_h, o_{h+1})\), and set \( N_h(a, \bar{a}) \leftarrow N_h(a, \bar{a}) + e_{o_h} e_{o_{h-1}}^T \).

11: set \( M_h(o_h, a, \bar{a}) \leftarrow M_h(o_h, a, \bar{a}) + e_{o_h} e_{o_{h-1}}^T \).

12: \( \mathcal{B}_h(a, \bar{a}) \leftarrow \bigcap_{o \in \mathcal{O}} \{ \theta \mid \| \mathcal{B}_h(a, o; \theta) N_h(a, \bar{a}) - M_h(o, a, \bar{a}) \|_F \leq \gamma_k \} \).

13: construct the confidence set \( \Theta_{k+1} = \bigcap_{(h, a, \bar{a}) \in [H-1] \times \mathcal{A}^2} \mathcal{B}_h(a, \bar{a}) \cap b \).

14: **Output:** \( \pi_k \) where \( k \) is sampled uniformly from \([K]\).

Confidence Bound (OOM-UCB, algorithm 1), is an optimistic algorithm which heavily exploits the OOM representation to obtain valid uncertainty estimates of the parameters of the underlying model.

To condense notation in Algorithm 1 we denote the parameters of a POMDP as \( \theta = (\mathcal{T}, \mathcal{O}, \mu_1) \). We denote \( V^\pi(\hat{\theta}) \) as the value of policy \( \pi \) if the underlying POMDP has parameter \( \theta \). We also write the parameters of the OOM \((b_0(\theta), \mathcal{B}_h(a, o; \theta))\) as a function of parameter \( \theta \), where the dependency is specified as in (1). We adopt the convention that at the 0-th step, the observation \( o_0 \) and state \( s_0 \) are always set to be some fixed dummy observation and state, and, starting from \( s_0 \), the environment transitions to \( s_1 \) with distribution \( \mu_1 \) regardless of what action \( a_0 \) is taken.

At a high level, Algorithm 1 is an iterative algorithm that, in each iteration, (a) computes an optimistic policy and model by maximizing the value (Line 4) subject to a given confidence set constraint, (b) collects data using the optimistic policy, and (c) incorporates the data into an updated confidence set for the OOM parameters (Line 5-13). The first two parts are straightforward, so we focus the discussion on computing the confidence set. We remark that in general the optimization in Line 4 may not be solved in polynomial time (see discussions of the computational complexity after Theorem 3).

First, since \( b_0 \) in (1) is simply the probability over observations at the first step, our confidence set for \( b_0 \) in Line 6 is simply based on counting the number of times each observation appears in the first step and Hoeffding’s concentration inequality.

Our construction of the confidence sets for the operators \( \{\mathcal{B}_h(a, o)\}_{h,a,o} \) is inspired by the method-of-moments estimator in HMM literature [Hsu et al. 2012]. Consider two fixed actions \( a, \bar{a}, \) and an arbitrary distribution over \( s_{h-1} \). Let \( P_h(a, \bar{a}), Q_h(o, a, \bar{a}) \in \mathbb{R}^{O \times O} \) be the probability matrices such that

\[
\begin{align*}
[P_h(a, \bar{a})]_{o', o''} &= \mathbb{P}(o_h = o', o_{h-1} = o'' | a_h = a, a_{h-1} = \bar{a}), \\
[Q_h(o, a, \bar{a})]_{o', o''} &= \mathbb{P}(o_{h+1} = o', o_h = o, o_{h-1} = o'' | a_h = a, a_{h-1} = \bar{a}).
\end{align*}
\]

It can be verified that \( \mathcal{B}_h(a, \bar{a}) P_h(a, \bar{a}) = \mathcal{Q}_h(a, o, \bar{a}) \) (Fact 17 in the appendix). Our confidence set construction (Line 12) in Algorithm 1 is based on this fact: we replace the probability matrices \( P, Q \) by empirical estimates \( N, M \), and we use concentration inequalities to determine the width of the confidence set. Finally, our overall confidence set for the parameters \( \theta \) is simply the intersection of the confidence sets for all induced operators and \( b_0 \), additionally incorporating the constraint on \( \sigma_{\min}(\hat{\Theta}_h) \) from Assumption 1.
3.3 Theoretical guarantees

Our OOM-UCB algorithm enjoys the following sample complexity guarantee.

**Theorem 3.** For any $\varepsilon \in (0, H]$, there exists $K_{\text{max}} = \text{poly}(H, S, A, O, \alpha^{-1})/\varepsilon^2$ and an absolute constant $c_1$, such that for any POMDP that satisfies Assumption 1 if we set hyperparameters $\beta_k = c_1 \sqrt{k \log(KAOH)}$, $\gamma_k = \sqrt{S \beta_k}/\alpha$, and $K \geq K_{\text{max}}$, then the output policy $\hat{\pi}$ of Algorithm 1 will be $\varepsilon$-optimal with probability at least $2/3$.

Theorem 3 claims that in polynomially many iterations of the outer loop, Algorithm 1 learns a near-optimal policy for any undercomplete POMDP that satisfies Assumption 1. Since our algorithm only uses $O(H^2 A^2)$ samples per iteration of the outer loop, this implies that the sample complexity is also $\text{poly}(H, S, A, O, \alpha^{-1})/\varepsilon^2$. We remark that the $1/\varepsilon^2$ dependence is optimal, which follows from standard concentration arguments. To the best of our knowledge, this is the first sample efficiency result for learning a class of POMDPs where exploration is essential.

While Theorem 3 does guarantee sample efficiency, Algorithm 1 is not computationally efficient due to that the computation of the optimistic policy (Line 4) may not admit a polynomial time implementation, which should be expected given the aforementioned computationally complexity results. We now turn to a further restricted (and interesting) subclass of POMDPs where we can address both the computational and statistical challenges.

4 Results for POMDPs with Deterministic Transition.

In this section, we complement our main result by investigating the class of POMDPs with deterministic transitions, where both computational and statistical efficiency can be achieved. We say a POMDP is of deterministic transition if both its transition and initial distribution are deterministic, i.e, if the entries of matrices $\{\mathbb{T}_h\}_h$ and vector $\mu_1$ are either 0 or 1. We remark that while deterministic dynamics avoids computational barriers, it does not mitigate the need for exploration.

Instead of Assumption 1 for the deterministic transition case, we require that the columns of the observation matrices $\mathbb{O}_h$ are well-separated.

**Assumption 2.** For any $h \in [H]$, $\min_{s \neq s'} \|\mathbb{O}_h(\cdot|s) - \mathbb{O}_h(\cdot|s')\| \geq \xi$.

Assumption 2 guarantees that observation distributions for different states are sufficiently different, by at least $\xi$ in Euclidean norm. It does not require that the POMDP is undercomplete, and, in fact, is strictly weaker than Assumption 1. In particular, for undercomplete models, $\min_{s \neq s'} \|\mathbb{O}_h(\cdot|s) - \mathbb{O}_h(\cdot|s')\| \geq \sqrt{2} \sigma_{\min}(\mathbb{O}_h)$, and so Assumption 1 implies Assumption 2 for $\xi = \sqrt{2} \alpha$.

Leveraging deterministic transitions, we can design a specialized algorithm (Algorithm 2 in the appendix) that learns an $\varepsilon$-optimal policy using polynomially many samples and in polynomial time. We present the formal theorem here, and refer readers to Appendix D for more details.

**Theorem 4.** For any $p \in (0, 1]$, there exists an algorithm such that for any POMDP with deterministic transitions that satisfies Assumption 2 within $O(H^2 SA \log(HSA/p)/(\min\{\varepsilon/\sqrt{OH}, \xi\})^2)$ samples and computations, the output policy of the algorithm is $\varepsilon$-optimal with probability at least $1 - p$.

5 Analysis Overview

In this section, we provide an overview of the proof of our main result—Theorem 3. Please refer to Appendix C for the full proof.
We start our analysis by noticing that the output policy \( \hat{\pi} \) of Algorithm 1 is uniformly sampled from \( \{\pi_k\}_{k=1}^K \) computed in the algorithm. If we can show that
\[
(1/K) \sum_{k=1}^K V^* - V^{\pi_k} \leq \varepsilon/10,
\] (4)
then at least a 2/3 fraction of the policies in \( \{\pi_k\}_{k=1}^K \) must be \( \varepsilon \)-optimal, and uniform sampling would find such a policy with probability at least 2/3. Therefore, our proof focuses on achieving (4).

We begin by conditioning on the event that for each iteration \( k \), our constructed confidence set \( \Theta_k \) in fact contains the true parameters \( \theta^* = (T, \emptyset, \mu) \) of the POMDP. This holds with high probability and is achieved by setting the widths \( \beta_k \) and \( \gamma_k \) appropriately (see Lemma 14 in the appendix).

### 5.1 Bounding suboptimality in value by error in density estimation.

Line 4 of Algorithm 1 computes the greedy policy \( \pi_k \leftarrow \arg\max_{\pi} \max_{\hat{\theta} \in \Theta_k} V^\pi(\hat{\theta}) \) with respect to the current confidence set \( \Theta_k \). Let \( \theta_k \) denote the maximizing model parameters in the \( k \)-th iteration. As \( (\pi_k, \theta_k) \) are optimistic, we have \( V^* \equiv V^*(\theta^*) \leq V^{\pi_k}(\theta_k) \) for all \( k \in [K] \). Thus, for any \( k \in [K] \):
\[
V^* - V^{\pi_k}(\theta_k) \leq H \sum_{o_H, \ldots, o_1} |P^\pi_{\theta_k}(o_H, \ldots, o_1) - P^\pi(\theta)(o_H, \ldots, o_1)|,
\]
where \( P^\pi_{\theta} \) denotes the probability measure over observations under policy \( \pi \) for POMDP with parameters \( \theta \). The second inequality holds because the cumulative reward is a function of observations \( (o_H, \ldots, o_1) \) and is upper bounded by \( H \). This upper bounds the suboptimality in value by the total variation distance between the \( H \)-step observation distributions.

Next, note that we can always choose the greedy policy \( \pi_k \) to be deterministic, i.e., the probability to take any action given a history is either 0 or 1. This allows us to define the following set for any deterministic policy \( \pi \):
\[
\Gamma(\pi, H) := \{\tau_H = (o_H, \ldots, a_1, o_1) \mid \pi(a_{H-1}, \ldots, a_1|o_H, \ldots, o_1) = 1\}.
\]
In words, \( \Gamma(\pi, H) \) is a set of all the observation and action sequences of length \( H \) that could occur under the \( \pi \). For any policy \( \pi \), there is a one-to-one correspondence between \( \Theta^H \) and \( \Gamma(\pi, H) \) and moreover, for any sequence \( \tau_H = (o_H, \ldots, a_1, o_1) \in \Gamma(\pi, H) \), we have:
\[
p(\tau_H; \theta) := P_\theta(o_H, \ldots, o_1|a_{H-1}, \ldots, a_1) = P^\pi_{\theta}(o_H, \ldots, o_1).
\]
Combining this with (5) and summing over all episodes, we conclude that:
\[
\sum_{k=1}^K (V^* - V^{\pi_k}) \leq H \sum_{k=1}^K \sum_{\tau_H \in \Gamma(\pi_k, H)} |p(\tau_H; \theta_k) - p(\tau_H; \theta^*)|.
\]
This upper bounds the suboptimality in value by errors in estimating the conditional probabilities.

### 5.2 Bounding error in density estimation by error in estimating operators.

For the next step, we leverage the OOM representation to bound the difference between the conditional probabilities \( p(\tau_H; \theta_k) \) and \( p(\tau_H; \theta^*) \). Recall that from (2), the conditional probability can be written as a product of the observable operators for each step and \( b_0 \). Therefore, for any two parameters \( \hat{\theta} \) and \( \theta \), we have following relation for any sequence \( \tau_H = (o_H, \ldots, a_1, o_1) \):
\[
p(\tau_H; \hat{\theta}) - p(\tau_H; \theta) = e_{a_H}^T \cdot B_{H-1}(a_{H-1}, o_{H-1}; \hat{\theta}) \cdots B_1(a_1, o_1; \hat{\theta}) \cdots b_0(\hat{\theta}) - b_0(\theta)
\]
\[
+ \sum_{h=1}^{H-1} e_{a_H}^T \cdot B_{H-1}(a_{H-1}, o_{H-1}; \hat{\theta}) \cdots [B_h(a_h, o_h; \hat{\theta}) - B_h(a_h, o_h; \theta)] \cdots B_1(a_1, o_1; \theta) \cdot b_0(\theta).
\]
This relates the difference \( p(\tau_H; \hat{\theta}) - p(\tau_H; \theta) \) to the differences in operators and \( b_0 \). Formally, with further relaxation and summation over all sequence in \( \Gamma(\pi, H) \), we have the following lemma (also see Lemma 10 in Appendix C).

**Lemma 5.** Given a deterministic policy \( \pi \) and two sets of undercomplete POMDP parameters \( \theta = (\Theta, \mathbb{T}, \mu_1) \) and \( \hat{\theta} = (\hat{\Theta}, \hat{\mathbb{T}}, \hat{\mu}_1) \) with \( \sigma_{\min}(\hat{\Theta}) \geq \alpha \), we have

\[
\begin{align*}
&\sum_{\tau_H \in \Gamma(\pi, H)} |p(\tau_H; \hat{\theta}) - p(\tau_H; \theta)| \leq \frac{\sqrt{S}}{\alpha} \left( \|b_0(\hat{\theta}) - b_0(\theta)\|_1 + \sum_{(a, o) \in \mathcal{A} \times \mathcal{O}} \|B_1(a, o; \hat{\theta}) - B_1(a, o; \theta)\|b_0(\theta)\|_1 \\
&\quad + \sum_{h=2}^{H-1} \sum_{(a, \hat{a}, o) \in \mathcal{A}^2 \times \mathcal{O}} \sum_{s=1}^{S} \left\| (B_h(a, o; \hat{\theta}) - B_h(a, o; \theta)) \Theta_h T_{h-1}(\hat{\alpha}) e_s \right\| \left\| \mathbb{P}_\theta(s_{h-1} = s) \right\|_1 \right).
\end{align*}
\]

(6)

This lemma suggests that if we could estimate the operators accurately, we would have small value sub-optimality. However, Assumption 1 is not sufficient for parameter recovery. It is possible that in some step \( h \), there exists a state \( s_h \) that can be reached with only very small probability no matter what policy is used. Since we cannot collect many samples from \( s_h \), it is not possible to estimate the corresponding component in the operator \( B_h \). In other words, we cannot hope to make \( \|B_h(a, o; \hat{\theta}) - B_h(a, o; \theta)\|_1 \) small in our setting.

To proceed, it is crucial to observe that the third term on the RHS of (6), in fact the operator error \( B_h(a, o; \hat{\theta}) - B_h(a, o; \theta) \) projected onto the direction \( \Theta_h T_{h-1}(\hat{\alpha}) e_s \) and additionally reweighted by the probability of visiting state \( s \) in step \( h - 1 \). Therefore, if \( s \) is hard to reach, the weighting probability will be very small, which means that even though we cannot estimate \( B_h(a, o; \theta) \) accurately in the corresponding direction, it has a negligible contribution to the density estimation error (LHS of (6)).

### 5.3 Bounding error in estimating operators by OOM-UCB algorithm

By Lemma 5, we only need to bound the error in operators reweighted by visitation probability. This is achieved by a careful design of the confidence sets in the OOM-UCB algorithm. This construction is based on the method of moments, which heavily exploits the undercompleteness of the POMDP. To showcase the main idea, we focus on bounding the third term on the RHS of (6).

Consider a fixed \( (o, a, \tilde{a}) \) tuple, a fixed step \( h \in [H] \), and a fixed iteration \( k \in [K] \). We define moment matrices \( P_h(a, \tilde{a}), Q_h(a, a, \tilde{a}) \in \mathbb{R}^{O \times O} \) as in (5) for distribution on \( s_{h-1} \) that equals \( (1/k) \cdot \sum_{t=1}^{k} \mathbb{P}_{\theta^t}(s_{h-1} = \cdot) \). We also denote \( \hat{P}_h(a, \tilde{a}) = N_h(a, \tilde{a})/k, \hat{Q}_h(a, a, \tilde{a}) = M_h(a, a, \tilde{a})/k \) for \( N_h, M_h \) matrices after the update in the \( k \)-th iteration of Algorithm 1. By martingale concentration, it is not hard to show that with high probability:

\[
\|P_h(a, \tilde{a}) - \hat{P}_h(a, \tilde{a})\|_F \leq \hat{O}(1/\sqrt{k}), \quad \|Q_h(a, a, \tilde{a}) - \hat{Q}_h(a, a, \tilde{a})\|_F \leq \hat{O}(1/\sqrt{k}).
\]

Additionally, we can show that for the true operator and the true moments, we have \( B_h(a, o; \theta^*) P_h(a, \tilde{a}) = Q_h(a, o, \tilde{a}) \). Meanwhile, by the construction of our confidence set \( \Theta_{k+1} \), we know that for any \( \hat{\theta} \in \Theta_{k+1} \), we have\[
\|B_h(a, o; \hat{\theta}) \hat{P}_h(a, \tilde{a}) - \hat{Q}_h(a, a, \tilde{a})\|_F \leq \gamma_k/k.
\]

Combining all relations above, we see that \( B_h(a, o; \hat{\theta}) \) is accurate in the directions spanned by \( P_h(a, \tilde{a}) \), which, by definition, are directions frequently visited by the previous policies \( \{\pi_t\}_{t=1}^{k} \). Formally, we have the following lemma (also see Lemma 15 in Appendix C), which allows us to further bound the third term on the RHS of (6) using the algebraic transformation in Lemma 16.

\[\text{Formally, we have the following lemma (also see Lemma 15 in Appendix C), which allows us to further bound the third term on the RHS of (6) using the algebraic transformation in Lemma 16.}\]
Lemma 6. With probability at least $1 - \delta$, for all $k \in [K]$, for any $\hat{\theta} = (\hat{\Theta}, \hat{T}, \hat{\mu}_1) \in \Theta_{k+1}$ and $(o, a, \tilde{a}, h) \in \mathcal{O} \times \mathcal{A}^2 \times \{2, \ldots, H-1\}$, and $t = \log(KAOh/\delta)$, we have

$$\sum_{s=1}^{S} \left\| \left( B_h(a, o; \hat{\theta}) - B_h(a, o; \theta^*) \right) \right\|_1 \sum_{t=1}^{k} \mathbb{P}_{\hat{\theta}}(s_{h-1} = s) \leq O\left( \sqrt{\frac{kS^2Oh}{\alpha^4}} \right).$$

6 Conclusion

In this paper, we give a sample efficient algorithm for reinforcement learning in undercomplete POMDPs. Our results leverage a connection to the observable operator model and employ a refined error analysis. To our knowledge, this gives the first provably efficient algorithm for strategic exploration in partially observable environments.

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A Notation

Below, we introduce some notations that will be used in appendices.

| notation | definition |
|----------|------------|
| \( n^k \) | value of \( n \) after the update in the \( k^{th} \) iteration of Algorithm I |
| \( N_h^k(a, \tilde{a}) \) | value of \( N_h(a, \tilde{a}) \) after the update in the \( k^{th} \) iteration of Algorithm I |
| \( M_h^k(o, a, \tilde{a}) \) | value of \( M_h(o, a, \tilde{a}) \) after the update in the \( k^{th} \) iteration of Algorithm I |
| \( \theta \) | a parameter triple \( (T, O, \mu_1) \) of a POMDP |
| \( \theta^* \) | the groundtruth POMDP parameter triple |
| \( \text{POMDP}(\theta) \) | POMDP\((H, \mathcal{A}, \mathcal{O}, T, \mathcal{O}, r, \mu_1)\) |
| \( \tau_h \) | a length-\( h \) trajectory: \( \tau_h = [a_h, o_h, \ldots, a_1, o_1] \in (\mathcal{A} \times \mathcal{O})^h \) |
| \( \Gamma(\pi, h) \) | \( \{\tau_h = (a_h, o_h, \ldots, a_1, o_1) \mid \pi(a_h, \ldots, a_1| o_h, \ldots, o_1) = 1\} \) |
| \( b(\tau_h; \theta) \) | \( B_h(a_h, o_h; \theta) \cdots B_1(a_1, o_1; \theta) \cdot b_0(\theta) \) |
| \( \mathbb{P}_\pi(s_h = s) \) | probability of visiting state \( s \) at \( h^{th} \) step when executing policy \( \pi \) on POMDP(\( \theta \)) |
| \( 1(x = y) \) | equal to 1 if \( x = y \) and 0 otherwise. |
| \( e_o \) | an \( O \)-dimensional vector with \( (e_o)_i = 1(o = i) \) |
| \( (X)_o \) | the \( o^{th} \) column of matrix \( X \) |
| \( I_n \) | \( n \times n \) identity matrix |
| \( C_{\text{poly}} \) | \( \text{poly}(S, O, A, H, 1/\alpha, \log(1/\delta)) \) |
| \( \ell \) | \( \log(AOHK/\delta) \) |

Let \( x \in \mathbb{R}^{n_x}, y \in \mathbb{R}^{n_y} \) and \( z \in \mathbb{R}^{n_z} \). We denote by \( x \otimes y \otimes z \) the tensor product of vectors \( x, y \) and \( z \), an \( n_x \times n_y \times n_z \) tensor with \((i, j, k)\)th entry equal to \( x_i y_j z_k \). Let \( X \in \mathbb{R}^{n_X \times m}, Y \in \mathbb{R}^{n_Y \times m} \) and \( Z \in \mathbb{R}^{n_Z \times m} \). We generalize the notation of tensor product to matrices by defining \( X \otimes Y \otimes Z = \sum_{l=1}^{m} (X)_l \otimes (Y)_l \otimes (Z)_l \), which is an \( n_X \times n_Y \times n_Z \) tensor with \((i, j, k)\)th entry equal to \( \sum_{l=1}^{m} x_{il} y_{jl} z_{kl} \).

Let \( X \) be a random variable taking value in \([m]\), we denote by \( \mathbb{P}(X = \cdot) \) an \( m \)-dimensional vector whose \( i^{th} \) entry is \( \mathbb{P}(X = i) \).

B Proof of Hardness Results

The hard examples constructed below are variants of the ones used in [Krishnamurthy et al. (2016)].

**Proposition 1.** For any algorithm \( \mathcal{A} \), there exists an overcomplete POMDP \((S > O)\) with \( S \) and \( O \) being small constants, which satisfies \( \sigma_{\min}(\mathbb{O}_h) = 1 \) for all \( h \in [H] \), such that algorithm \( \mathcal{A} \) requires at least \( \Omega(A^{H-1}) \) samples to ensure learning a \((1/4)\)-optimal policy with probability at least 1/2.

\(^3\)Note that this definition is different from the one used in Section 5 where \( \tau_h = [o_h, \ldots, a_1, o_1] \in \mathcal{O} \times (\mathcal{A} \times \mathcal{O})^{h-1} \) does not include the action \( a_h \) at \( h^{th} \) step.

\(^4\)WLOG, all policies considered in this paper are deterministic. Also note that the trajectory in \( \Gamma(\pi, h) \) contains \( a_h \), which is different from the definition in Section 5.
We have the following key observations:

1. **STATE** There are four states: two good states $g_1$ and $g_2$ and two bad states $b_1$ and $b_2$. The initial state is picked uniformly at random.

2. **OBSERVATION** There are only two different observations $u_1$ and $u_2$. At step $h \in [H-1]$, we always observe $u_1$ at $g_1$ and $b_1$, and observe $u_2$ at $g_2$ and $b_2$. At step $H$, we always observe $u_1$ at good states and $u_2$ at bad states. We can immediately verify that $\sigma_{\text{min}}(O_h) = 1$ for all $h \in [H]$.

3. **REWARD** There is no reward in the first $H-1$ steps (i.e. $r_h = 0$ for all $h \in [H-1]$). At step $H$, we receive reward 1 if we observe $u_1$ and no reward otherwise (i.e. $r_H(o) = 1(o = u_1)$).

4. **TRANSITION** There is one good action $a^*_h$ and $A−1$ bad actions for each $h \in [H-1]$. At step $h \in [H-1]$, suppose we are at a good state ($g_1$ or $g_2$), then we will transit to $g_1$ or $g_2$ uniformly at random if we take $a^*_h$. Otherwise, we transit to $b_1$ or $b_2$ uniformly at random. In contrast, if we are at a bad state ($b_1$ or $b_2$), we will always transit to $b_1$ or $b_2$ uniformly at random no matter what action we take. Note that two good (bad) states are equivalent in terms of transition.

We have the following key observations:

1. Once we are at bad states, we always stay at bad states.

2. We have
\[
P(o_{1:H−1} = z \mid a_{1:H−1}, o_H) = \frac{1}{2^{H−1}}
\]
for any $z \in \{u_1, u_2\}^{H−1}$ and $(a_{1:H−1}, o_H) \in [A]^{H−1} \times \{u_1, u_2\}$

Therefore, the observations at the first $H−1$ steps provide no information about the underlying transition. The only useful information is the last observation $o_H$ which tells us whether we end in good states or not.

3. The optimal policy is unique and is to execute the good action sequence $(a^*_1, \ldots, a^*_{H−1})$ regardless of the observations.

Based on the observations above, this is equivalent to a multi-arm bandits problem with $A^{H−1}$ arms. Therefore, we cannot do better than Brute-force search, which has sample complexity at least $\Omega(A^{H−1})$.

**Proposition 2.** For any algorithm $\mathfrak{A}$, there exists an undercomplete POMDP ($S \leq O$) with $S$ and $O$ being small constants, such that algorithm $\mathfrak{A}$ requires at least $\Omega(A^{H−1})$ samples to ensure learning a $(1/4)$-optimal policy with probability at least $1/2$.

**Proof.** We continue to use the POMDP constructed in Proposition 1 and slightly modify it by splitting $u_2$ into another 4 different observations $\{q_1, q_2, q_3, q_4\}$, so in the new POMDP ($O = 5 > S = 4$), we will observe a $q_i$ picked uniformly at random from $\{q_1, q_2, q_3, q_4\}$ when we are 'supposed' to observe $u_2$. It is easy to see the modification does not change its hardness.

**C Analysis of OMM-UCB**

Throughout the proof, we use $\tau_h$ to denote a length-$h$ trajectory: $[a_h, o_h, \ldots, a_1, o_1] \in (A \times \Theta)^h$. Note that this definition is different from the one used in Section 5, where $\tau_h = [o_h, \ldots, a_1, o_1] \in \Theta \times (A \times \Theta)^{h−1}$ does not include the action $a_h$ at $h$th step. Besides, we define
\[
\Gamma(\pi, h) = \{\tau_h = (a_h, o_h, \ldots, a_1, o_1) \mid \pi(a_h, \ldots, a_1 | o_h, \ldots, o_1) = 1\},
\]
which is also different from the definition in Section [5] whether \( a_h \) is not included.

Please refer to Appendix [A] for definitions of frequently used notations.

C.1 Bounding the error in belief states

In this subsection, we will bound the error in (unnormalized) belief states, i.e., \( b(\tau_h; \theta) - b(\tau_h; \hat{\theta}) \) by the error in operators reweighed by the probability distribution of visited states.

We start by proving the following lemma that helps us decompose the error in belief states inductively.

**Lemma 7.** Given a deterministic policy \( \pi \) and two set of POMDP parameters \( \hat{\theta} = (\hat{\Theta}, \hat{T}, \hat{\mu}_1) \) and \( \theta = (\Theta, T, \mu_1) \), for all \( h \geq 1 \) and \( X \in \{ I_O, \hat{O}_{h+1} \} \), we have

\[
\sum_{\tau_h \in \Gamma(\pi, h)} \left\| X \left( b(\tau_h; \theta) - b(\tau_h; \hat{\theta}) \right) \right\|_1 \leq \sum_{\tau_{h-1} \in \Gamma(\pi, h-1)} \left\| \hat{O}_{h}^\dagger \left( b(\tau_{h-1}; \theta) - b(\tau_{h-1}; \hat{\theta}) \right) \right\|_1 + \sum_{\tau_h \in \Gamma(\pi, h)} \left\| X \left( B_h(a_h, o_h; \hat{\theta}) - B_h(a_h, o_h; \theta) \right) b(\tau_{h-1}; \theta) \right\|_1.
\]

**Proof.** By the definition of \( b(\tau_h; \theta) \) and \( b(\tau_h; \hat{\theta}) \),

\[
\sum_{\tau_h \in \Gamma(\pi, h)} \left\| X \left( b(\tau_h; \theta) - b(\tau_h; \hat{\theta}) \right) \right\|_1 = \sum_{\tau_h \in \Gamma(\pi, h)} \left\| X \left( B_h(a_h, o_h; \theta) b(\tau_{h-1}; \theta) - B_h(a_h, o_h; \hat{\theta}) b(\tau_{h-1}; \hat{\theta}) \right) \right\|_1 \leq \sum_{\tau_h \in \Gamma(\pi, h)} \left\| X B_h(a_h, o_h; \hat{\theta}) \left( b(\tau_{h-1}; \theta) - b(\tau_{h-1}; \hat{\theta}) \right) \right\|_1 + \sum_{\tau_h \in \Gamma(\pi, h)} \left\| X \left( B_h(a_h, o_h; \hat{\theta}) - B_h(a_h, o_h; \theta) \right) b(\tau_{h-1}; \theta) \right\|_1.
\]

The first term can be bounded as following,

\[
\sum_{\tau_h \in \Gamma(\pi, h)} \left\| X B_h(a_h, o_h; \hat{\theta}) (b(\tau_{h-1}; \theta) - b(\tau_{h-1}; \hat{\theta})) \right\|_1 = \sum_{\tau_h \in \Gamma(\pi, h)} \left\| X \hat{T}_{h+1} \hat{\Theta}_h(a_h) \text{diag}(\hat{\Theta}_h(o_h | \cdot)) \hat{O}_{h}^\dagger \left( b(\tau_{h-1}; \theta) - b(\tau_{h-1}; \hat{\theta}) \right) \right\|_1 \leq \sum_{\tau_h \in \Gamma(\pi, h)} \sum_i \left\| \left( X \hat{T}_{h+1} \hat{\Theta}_h(a_h) \text{diag}(\hat{\Theta}_h(o_h | \cdot)) \right) \hat{O}_{h}^\dagger \left( b(\tau_{h-1}; \theta) - b(\tau_{h-1}; \hat{\theta}) \right) \right\|_1 \leq \sum_{\tau_h \in \Gamma(\pi, h)} \sum_i \hat{\Theta}_h(o_h | i) \left\| \hat{O}_{h}^\dagger \left( b(\tau_{h-1}; \theta) - b(\tau_{h-1}; \hat{\theta}) \right) \right\|_1 = \sum_{\tau_h \in \Gamma(\pi, h)} \sum_i \hat{O}_{h}^\dagger \left( b(\tau_{h-1}; \theta) - b(\tau_{h-1}; \hat{\theta}) \right) _i,
\]

where the inequality is by triangle inequality, and the last identity follows from \( \hat{T}_{h}(a_h) \) (when \( X = \hat{O}_{h+1}^\dagger \)) and \( \hat{O}_{h+1} \hat{T}_{h}(a_h) \) (when \( X = I_O \)) having columns with \( \ell_1 \)-norm equal to 1.
As $\pi$ is deterministic, $a_h$ is unique given $\tau_{h-1}$ and $o_h$. Therefore,

$$
\sum_{\tau_h \in \Gamma(\pi, h)} \sum_i \hat{\mathbf{O}}_h(o_h \mid i) \left\| \left( \hat{\mathbf{O}}_h^\dagger \left( \mathbf{b}(\tau_{h-1}; \theta) - \mathbf{b}(\tau_{h-1}; \hat{\theta}) \right) \right)_i \right\|
$$

$$
= \sum_{\tau_{h-1} \in \Gamma(\pi, h-1)} \sum_{\tau_h \in \Gamma(\pi, h)} \sum_i \hat{\mathbf{O}}_h(o_h \mid i) \left\| \left( \hat{\mathbf{O}}_h^\dagger \left( \mathbf{b}(\tau_{h-1}; \theta) - \mathbf{b}(\tau_{h-1}; \hat{\theta}) \right) \right)_i \right\|
$$

$$
= \sum_{\tau_{h-1} \in \Gamma(\pi, h-1)} \sum_{\tau_h \in \Gamma(\pi, h)} \sum_i \hat{\mathbf{O}}_h(o_h \mid i) \left\| \left( \hat{\mathbf{O}}_h^\dagger \left( \mathbf{b}(\tau_{h-1}; \theta) - \mathbf{b}(\tau_{h-1}; \hat{\theta}) \right) \right)_i \right\|
$$

$$
= \sum_{\tau_{h-1} \in \Gamma(\pi, h-1)} \sum_i \left\| \left( \hat{\mathbf{O}}_h^\dagger \left( \mathbf{b}(\tau_{h-1}; \theta) - \mathbf{b}(\tau_{h-1}; \hat{\theta}) \right) \right)_i \right\|_1.
$$

which completes the proof. $\square$

By applying Lemma 7 inductively, we can bound the error in belief states by the projection of errors in operators on preceding belief states.

**Lemma 8.** Given a deterministic policy $\pi$ and two sets of undercomplete POMDP parameters $\theta = (\mathbf{O}, \mathbb{T}, \mu_1)$ and $\hat{\theta} = (\hat{\mathbf{O}}, \hat{\mathbb{T}}, \hat{\mu}_1)$ with $\sigma_{\text{min}}(\hat{\mathbf{O}}) \geq \alpha$, for all $h \geq 1$, we have

$$
\sum_{\tau_h \in \Gamma(\pi, h)} \left\| \mathbf{b}(\tau_h; \theta) - \mathbf{b}(\tau_h; \hat{\theta}) \right\|_1 
\leq \frac{\sqrt{S}}{\alpha} \sum_{j=1}^h \sum_{\tau_j \in \Gamma(\pi, j)} \left\| \left( \mathbf{B}_j(a_j, o_j; \theta) - \mathbf{B}_j(a_j, o_j; \hat{\theta}) \right) \mathbf{b}(\tau_{j-1}; \theta) \right\|_1 + \frac{\sqrt{S}}{\alpha} \left\| \mathbf{b}_0(\theta) - \mathbf{b}_0(\hat{\theta}) \right\|_1.
$$

**Proof.** Invoking Lemma 7 with $X = \hat{\mathbf{O}}_{j+1}^\dagger$, we have

$$
\sum_{\tau_j \in \Gamma(\pi, j)} \left\| \hat{\mathbf{O}}_{j+1}^\dagger \left( \mathbf{b}(\tau_{j}; \theta) - \mathbf{b}(\tau_{j}; \hat{\theta}) \right) \right\|_1 \leq \sum_{\tau_{j-1} \in \Gamma(\pi, j-1)} \left\| \hat{\mathbf{O}}_{j}^\dagger \left( \mathbf{b}(\tau_{j-1}; \theta) - \mathbf{b}(\tau_{j-1}; \hat{\theta}) \right) \right\|_1
$$

$$
+ \sum_{\tau_j \in \Gamma(\pi, j)} \left\| \hat{\mathbf{O}}_{j+1}^\dagger \left( \mathbf{b}(\tau_{j}; \theta) - \mathbf{b}(\tau_{j}; \hat{\theta}) \right) \right\|_1.
$$

(7)

Summing (7) over $j = 1, \ldots, h - 1$, we obtain

$$
\sum_{\tau_{h-1} \in \Gamma(\pi, h-1)} \left\| \hat{\mathbf{O}}_{h}^\dagger \left( \mathbf{b}(\tau_{h-1}; \theta) - \mathbf{b}(\tau_{h-1}; \hat{\theta}) \right) \right\|_1 \leq \sum_{j=1}^{h-1} \sum_{\tau_j \in \Gamma(\pi, j)} \left\| \hat{\mathbf{O}}_{j+1}^\dagger \left( \mathbf{b}(\tau_{j}; \theta) - \mathbf{b}(\tau_{j}; \hat{\theta}) \right) \right\|_1 + \left\| \hat{\mathbf{O}}_{1}^\dagger \left( \mathbf{b}_0(\theta) - \mathbf{b}_0(\hat{\theta}) \right) \right\|_1.
$$

(8)

Again, invoking Lemma 7 with $X = I_O$ gives

$$
\sum_{\tau_h \in \Gamma(\pi, h)} \left\| \mathbf{b}(\tau_h; \theta) - \mathbf{b}(\tau_h; \hat{\theta}) \right\|_1 \leq \sum_{\tau_{h-1} \in \Gamma(\pi, h-1)} \left\| \hat{\mathbf{O}}_{h}^\dagger \left( \mathbf{b}(\tau_{h-1}; \theta) - \mathbf{b}(\tau_{h-1}; \hat{\theta}) \right) \right\|_1
$$

$$
+ \sum_{\tau_h \in \Gamma(\pi, h)} \left\| \left( \mathbf{B}_h(a_h, o_h; \theta) - \mathbf{B}_h(a_h, o_h; \hat{\theta}) \right) \mathbf{b}(\tau_{h-1}; \theta) \right\|_1.
$$

(9)

Plugging (8) into (9), and using the fact that $\left\| \hat{\mathbf{O}}_{h}^\dagger \right\|_{1 \rightarrow 1} \leq \sqrt{S} \left\| \hat{\mathbf{O}}_{h}^\dagger \right\|_2 \leq \frac{\sqrt{S}}{\alpha}$ complete the proof. $\square$
The following lemma bounds the projection of any vector on belief states by its projection on the product of the observation matrix and the transition matrix, reweighted by the visitation probability of states.

**Lemma 9.** For any deterministic policy $\pi$, fixed $a_{h+1} \in \mathcal{A}$, $u \in \mathbb{R}^O$, and $h \geq 0$, we have

$$\sum_{a_{h+1} \in \mathcal{A}} \sum_{\tau_h \in \Gamma(\pi,h)} |u^T b([a_{h+1}, o_{h+1}, \tau_h]; \theta)| \leq \sum_{s=1}^{S} |u^T (\mathcal{O}_{h+2} \mathbb{T}_{h+1}(a_{h+1}))_s | \mathbb{P}_\theta(s_{h+1} = s).$$

**Proof.** By definition, for any $[a_{h+1}, o_{h+1}, \tau_h] \in \mathcal{A} \times \mathcal{O} \times \Gamma(\pi, h)$, we have

$$b([a_{h+1}, o_{h+1}, \tau_h]; \theta) = \mathcal{O}_{h+2} \mathbb{T}_{h+1}(a_{h+1}) \mathbb{P}_\theta(s_{h+1} = \cdot, [o_{h+1}, \tau_h]),$$

where $\mathbb{P}_\theta(s_{h+1} = \cdot, [o_{h+1}, \tau_h])$ is an $s$-dimensional vector, whose $i$th entry is equal to the probability of observing $[o_{h+1}, \tau_h]$ and reaching state $i$ at step $h + 1$ when executing policy $\pi$ in POMDP($\theta$).

Therefore,

$$\sum_{\tau_h \in \Gamma(\pi,h)} \sum_{o_{h+1} \in \mathcal{O}} |u^T b([a_{h+1}, o_{h+1}, \tau_h]; \theta)|$$

$$= \sum_{\tau_h \in \Gamma(\pi,h)} \sum_{o_{h+1} \in \mathcal{O}} |u^T \mathcal{O}_{h+2} \mathbb{T}_{h+1}(a_{h+1}) \mathbb{P}_\theta(s_{h+1} = \cdot, [o_{h+1}, \tau_h])|$$

$$\leq \sum_{\tau_h \in \Gamma(\pi,h)} \sum_{o_{h+1} \in \mathcal{O}} \sum_{s=1}^{S} |u^T \mathcal{O}_{h+2} \mathbb{T}_{h+1}(a_{h+1})_s | \mathbb{P}_\theta(s_{h+1} = s, [o_{h+1}, \tau_h])$$

$$= \sum_{s=1}^{S} |u^T \mathcal{O}_{h+2} \mathbb{T}_{h+1}(a_{h+1})_s | \left( \sum_{\tau_h \in \Gamma(\pi,h)} \sum_{o_{h+1} \in \mathcal{O}} \mathbb{P}_\theta(s_{h+1} = s, [o_{h+1}, \tau_h]) \right)$$

$$= \sum_{s=1}^{S} |u^T \mathcal{O}_{h+2} \mathbb{T}_{h+1}(a_{h+1})_s | \mathbb{P}_\theta(s_{h+1} = s). \quad \square$$

Combining Lemma 8 and Lemma 9 we obtain the target bound.

**Lemma 10.** Given a deterministic policy $\pi$ and two sets of undercomplete POMDP parameters $\theta = (\mathcal{O}, \mathbb{T}, \mu_1)$ and $\hat{\theta} = (\hat{\mathcal{O}}, \hat{\mathbb{T}}, \hat{\mu}_1)$ with $\sigma_{\min}(\hat{\mathcal{O}}) \geq \alpha$, for all $h \geq 1$, we have

$$\sum_{\tau_h \in \Gamma(\pi,h)} \| b(\tau_h; \theta) - b(\tau_h; \hat{\theta}) \|_1$$

$$\leq \frac{\sqrt{S}}{\alpha} \left\| b_0(\theta) - b_0(\hat{\theta}) \right\|_1 + \frac{\sqrt{S}}{\alpha} \sum_{(a, o) \in \mathcal{A} \times \mathcal{O}} \left\| \left( \mathbb{B}_1(a, o; \theta) - \mathbb{B}_1(a, o; \hat{\theta}) \right) b_0(\theta) \right\|_1$$

$$+ \frac{\sqrt{S}}{\alpha} \sum_{j=2}^{h} \sum_{(a, o, \hat{o}) \in \mathcal{A} \times \mathcal{O} \times \mathcal{O}} \left\| \left( \mathbb{B}_j(a, \hat{o}; \hat{\theta}) - \mathbb{B}_j(a, o; \theta) \right) \left( \mathcal{O}_j \mathbb{T}_{j-1}(\hat{\theta}) \right)_s \right\|_1 \mathbb{P}_\theta(s_{j-1} = s).$$

**Proof.** By Lemma 8

$$\sum_{\tau_h \in \Gamma(\pi,h)} \| b(\tau_h; \theta) - b(\tau_h; \hat{\theta}) \|_1$$

$$\leq \frac{\sqrt{S}}{\alpha} \sum_{j=2}^{h} \sum_{\tau_j \in \Gamma(\pi,j)} \left\| \left( \mathbb{B}_j(a_j, o_j; \theta) - \mathbb{B}_j(a_j, o_j; \hat{\theta}) \right) b(\tau_{j-1}; \theta) \right\|_1$$

$$+ \frac{\sqrt{S}}{\alpha} \sum_{\tau_1 \in \Gamma(\pi,1)} \left\| \left( \mathbb{B}_1(a_1, o_1; \theta) - \mathbb{B}_1(a_1, o_1; \hat{\theta}) \right) b_0(\theta) \right\|_1 + \frac{\sqrt{S}}{\alpha} \left\| b_0(\theta) - b_0(\hat{\theta}) \right\|_1. \quad (10)$$
Bounding the first term: note that $\Gamma(\pi, j) \subseteq \Gamma(\pi, j - 2) \times (\Theta \times \mathcal{A})^2$, so we have
\[
\sum_{\tau_j \in \Gamma(\pi, j)} \| (B_j(a_j, o_j; \hat{\theta}) - B_j(a_j, o_j; \theta)) b(\tau_j; \theta) \|_1 \\
\leq \sum_{\tau_{j-2} \in \Gamma(\pi, j - 2)} \sum_{a_{j-1} \in \Theta} \sum_{a_{j-1} \in \mathcal{A}} \sum_{a_j \in \Theta} \sum_{a_j \in \mathcal{A}} \| (B_j(a_j, o_j; \hat{\theta}) - B_j(a_j, o_j; \theta)) b([a_{j-1}, o_{j-1}, \tau_{j-2}]; \theta) \|_1 \\
= \sum_{(a_j, a_{j-1}, o_j) \in \mathcal{A}^2 \times \Theta} \sum_{\tau_{j-2} \in \Gamma(\pi, j - 2)} \sum_{a_{j-1} \in \Theta} \| (B_j(a_j, o_j; \hat{\theta}) - B_j(a_j, o_j; \theta)) b([a_{j-1}, o_{j-1}, \tau_{j-2}]; \theta) \|_1.
\]

(11)

We can bound (\circ) by Lemma 9 and obtain,
\[
\sum_{\tau_j \in \Gamma(\pi, j)} \| (B_j(a_j, o_j; \hat{\theta}) - B_j(a_j, o_j; \theta)) b(\tau_j; \theta) \|_1 \\
\leq \sum_{(a_{j-1}, o_j) \in \mathcal{A}^2 \times \Theta} S \| (B_j(a_j, o_j; \hat{\theta}) - B_j(a_j, o_j; \theta)) (\tilde{O}_j \tilde{T}_{j-1}(a_{j-1}))_s \|_1 \tilde{P}_\Theta(s_{j-1} = s) \\
= \sum_{(a, \tilde{a}, o) \in \mathcal{A}^2 \times \Theta} S \| (B_j(a, o; \hat{\theta}) - B_j(a, o; \theta)) (\tilde{O}_j \tilde{T}_{j-1}(\tilde{a}))_s \|_1 \tilde{P}_\Theta(s_{j-1} = s),
\]

(12)

where the identity only changes the notations $(a_j, a_{j-1}, o_j) \rightarrow (a, \tilde{a}, o)$ to make the expression cleaner.

Bounding the second term: note that $\Gamma(\pi, 1) \subseteq \Theta \times \mathcal{A}$, we have
\[
\sum_{\tau_1 \in \Gamma(\pi, 1)} \| (B_1(a_1, o_1; \hat{\theta}) - B_1(a_1, o_1; \theta)) b_0(\theta) \|_1 \\
\leq \sum_{(a, o) \in \mathcal{A} \times \Theta} \| (B_1(a, o; \hat{\theta}) - B_1(a, o; \hat{\theta})) b_0(\theta) \|_1.
\]

(13)

Plugging (12) and (13) into (10) completes the proof.

C.2 A hammer for studying confidence sets

In this subsection, we develop a martingale concentration result, which forms the basis of analyzing confidence sets.

We start by giving the following basic fact about POMDP. The proof is just some basic algebraic calculation so we omit it here.

Fact 11. In POMDP($\theta$), suppose $s_{h-1}$ is sampled from $\mu_{h-1}$, fix $a_{h-1} \equiv \tilde{a}$, and $a_h \equiv a$. Then the joint distribution of $(o_{h+1}, o_h, a_{h-1})$ is
\[
\mathbb{P}(o_{h+1} = \cdot, o_h = \cdot, o_{h-1} = \cdot) = (\mathbb{O}_{h+1} T_h(a)) \otimes \mathbb{O}_h \otimes (\mathbb{O}_{h-1} \text{diag}(\mu_{h-1}) T_{h-1}(\tilde{a})^\top).
\]

By slicing the tensor, we can further obtain
\[
\left\{ \begin{array}{l}
\mathbb{P}(o_{h+1} = \cdot, o_h = \cdot, o_{h-1} = \cdot) = \mathbb{O}_{h-1} \mu_{h-1}, \\
\mathbb{P}(o_h = \cdot, o_{h-1} = \cdot) = \mathbb{O}_h T_{h-1}(\tilde{a}) \text{diag}(\mu_{h-1}) \mathbb{O}_{h-1}^\top, \\
\mathbb{P}(o_{h+1} = \cdot, o_h = a, o_{h-1} = \cdot) = \mathbb{O}_{h+1} T_h(a) \text{diag}(\mathbb{O}_h(a \mid \cdot)) T_{h-1}(\tilde{a}) \text{diag}(\mu_{h-1}) \mathbb{O}_{h-1}^\top.
\end{array} \right.
\]

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A simple implication of Fact 11 is that if we execute policy \( \pi \) from step 1 to step \( h - 2 \), take \( \tilde{a} \) and \( a \) at step \( h - 1 \) and \( h \) respectively, then the joint distribution of \( (o_{h+1}^t, o_h^t, o_{h-1}^t) \) is the same as above except for replacing \( \mu_{h-1} \) with \( \mathbb{P}_\theta^\pi(s_{h-1} = \cdot) \).

Suppose we are given a set of sequential data \( \{(o_{h+1}^t, o_h^t, o_{h-1}^t)\}_{t=1}^N \) generated from POMDP(\( \theta \)) in the following way: at time \( t \), execute policy \( \pi_t \) from step 1 to step \( h - 2 \), take action \( \tilde{a} \) at step \( h - 1 \), and action \( a \) at step \( h \) respectively, and observe \( (o_{h+1}^t, o_h^t, o_{h-1}^t) \). Here, we allow the policy \( \pi_t \) to be adversarial, in the sense that \( \pi_t \) can be chosen based on \( \{(\pi_t, o_{h+1}^t, o_h^t, o_{h-1}^t)\}_{t=1}^{t-1} \). Define \( \mu_{h-1}^{adv} = \frac{1}{N} \sum_{t=1}^N \mathbb{P}_\theta^\pi_t(s_{h-1} = \cdot) \). Based on Fact 11 we define the following probability vector, matrices and tensor,

\[
\begin{align*}
P_{h-1} & = \mathbb{O}_{h-1} \mu_{h-1}^{adv}, \\
P_{h,h-1} & = \mathbb{O}_h \mathbb{T}_{h-1}(\tilde{a}) \text{diag}(\mu_{h-1}^{adv}) \mathbb{O}_h^\top, \\
P_{h+1,h,h-1} & = (\mathbb{O}_{h+1} \mathbb{T}_h(a)) \circ \mathbb{O}_h \circ (\mathbb{O}_{h-1} \text{diag}(\mu_{h-1}^{adv}) \mathbb{T}_{h-1}(\tilde{a})^\top) \\
P_{h+1,o,h-1} & = \mathbb{O}_{h+1} \mathbb{T}_h(a) \circ (\mathbb{O}_h(o | \cdot)) \mathbb{T}_{h-1}(\tilde{a}) \text{diag}(\mu_{h-1}^{adv}) \mathbb{O}_h^\top, \quad o \in \mathcal{O}.
\end{align*}
\]

Accordingly, we define their empirical estimates as below

\[
\begin{align*}
\hat{P}_{h-1} & = \frac{1}{N} \sum_{t=1}^N e_{o_{h-1}^t}, \\
\hat{P}_{h,h-1} & = \frac{1}{N} \sum_{t=1}^N e_{o_{h-1}^t} \otimes e_{o_{h-1}^t}, \\
\hat{P}_{h+1,h,h-1} & = \frac{1}{N} \sum_{t=1}^N e_{o_{h+1}^t} \otimes e_{o_h^t} \otimes e_{o_{h-1}^t}, \\
\hat{P}_{h+1,o,h-1} & = \frac{1}{N} \sum_{t=1}^N e_{o_{h+1}^t} \otimes e_{o_h^t} \otimes 1(o_h^t = o), \quad o \in \mathcal{O}.
\end{align*}
\]

**Lemma 12.** There exists an absolute constant \( c_1 \), s.t. the following concentration bound holds with probability at least \( 1 - \delta \)

\[
\max_{o \in \mathcal{O}} \left\| \hat{P}_{h+1,h,h-1} - P_{h+1,h,h-1} \right\|_F, \left\| \hat{P}_{h,h-1} - P_{h,h-1} \right\|_F, \left\| \hat{P}_{h+1,o,h-1} - P_{h+1,o,h-1} \right\|_F, \left\| \hat{P}_{h-1} - P_{h-1} \right\|_2 \leq c_1 \sqrt{\frac{\log(N/\delta)}{N}}.
\]

**Proof.** Let \( \mathcal{F}_t \) be the \( \sigma \)-algebra generated by \( \{\{\pi_i\}_{i=1}^{t+1}, \{(a_{i+1}^t, o_h^t, o_{h-1}^t)\}_{i=1}^t\} \). \( \mathcal{F}_t \) is a filtration. Define

\[
X_t = e_{o_{h+1}^t} \otimes e_{o_h^t} \otimes e_{o_{h-1}^t} - (\mathbb{O}_{h+1} \mathbb{T}_h(a)) \circ \mathbb{O}_h \circ (\mathbb{O}_{h-1} \text{diag}(\mathbb{P}_\theta^\pi_t(s_{h-1} = \cdot)) \mathbb{T}_{h-1}(\tilde{a})^\top).
\]

We have \( X_t \in \mathcal{F}_t \) and \( \mathbb{E}[X_t | \mathcal{F}_{t-1}] = \mathbb{E}[X_t | \pi_t] = 0 \), where the second identity follows from Fact 11. Moreover,

\[
\|X_t\|_F \leq \|X_t\|_1 \leq \|e_{o_{h+1}^t} \otimes e_{o_h^t} \otimes e_{o_{h-1}^t}\|_1 + \|(\mathbb{O}_{h+1} \mathbb{T}_h(a)) \circ \mathbb{O}_h \circ (\mathbb{O}_{h-1} \text{diag}(\mathbb{P}_\theta^\pi_t(s_{h-1} = \cdot)) \mathbb{T}_{h-1}(\tilde{a})^\top)\|_1 = 2,
\]

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where $\| \cdot \|_1$ denotes the entry-wise $\ell_1$-norm of the tensor. By Azuma-Hoeffding’s inequality, with probability at least $1 - \delta/2$,

$$\| \hat{P}_{h+1,h,h-1} - P_{h+1,h,h-1} \|_F = \frac{1}{N} \sum_{t=1}^{N} (\delta_{h(t)} \otimes \delta_{o(t)} \otimes \delta_{o_{h-1}(t)} - (\mathbb{O}_{h+1} \mathbb{T}_h(a)) \otimes \mathbb{O}_h \otimes (\mathbb{O}_{h-1} \text{diag}(P_{o_{h-1} = 1}^{\sigma_{h-1}} \mathbb{T}_{h-1}(\tilde{a})^\top)) \|_F$$

where $c_1$ is some absolute constant. Similarly, we can show that with probability at least $1 - \delta/2$,

$$\| \hat{P}_{h,h-1} - P_{h,h-1} \|_F \leq c_1 \sqrt{\frac{\log(N/\delta)}{N}}$$

and

$$\| \hat{P}_{h-1} - P_{h-1} \|_F \leq c_1 \sqrt{\frac{\log(N/\delta)}{N}}.$$  

Using the fact $\| \hat{P}_{h+1,o,h-1} - P_{h+1,o,h-1} \|_F \leq \| \hat{P}_{h+1,h,h-1} - P_{h+1,h,h-1} \|_F$ completes the whole proof.

**C.3 Properties of confidence sets**

For convenience of discussion, we divide the constraints in $\Theta_k$ into three categories as follows:

**Type-0 constraint:**

$$\| k \cdot b_0(\hat{\theta}) - n^k \|_2 \leq \beta_k$$

**Type-I constraint:**

$$\| B_1(a, o; \hat{\theta})N_k^o(a, \tilde{a}) - M_k^t(o, a, \tilde{a}) \|_F \leq \gamma_k$$

where $M_k^t$ and $N_k^o$ are actually equivalent to $O$-dimensional counting vectors because there is no observation (or only a dummy observation) at step 0, which implies each of them has only one non-zero column. With slight abuse of notation, we use $M_k^t$ and $N_k^o$ to denote their non-zero columns in the following proof.

**Type-II constraint:** for $2 \leq h \leq H - 1$,

$$\| B_h(a, o; \hat{\theta})N_k^h(a, \tilde{a}) - M_k^h(o, a, \tilde{a}) \|_F \leq \gamma_k$$

Recalling the definition of $n^k(\theta)$, $N_k^o(a, \tilde{a})$ and $M_k^t(o, a, \tilde{a})$ and applying Lemma[12] we get the following concentration results.

**Corollary 13.** Let $\theta^* = (\mathbb{T}, \mathbb{O}, \mu_1)$. By applying Lemma[12] directly, with probability at least $1 - \delta$, for all $k \in [K]$ and $(o, a, \tilde{a}) \in \mathcal{O} \times \mathcal{O}^2$, we have

\[
\begin{align*}
\left\| \frac{1}{k} n^k - \mathbb{O}_1 \mu_1 \right\|_2 &\leq \mathcal{O}\left(\sqrt{\frac{\ell}{k}}\right), \\
\left\| \frac{1}{k} N_k^o(a, \tilde{a}) - \mathbb{O}_1 \mu_1 \right\|_2 &\leq \mathcal{O}\left(\sqrt{\frac{\ell}{k}}\right), \\
\left\| \frac{1}{k} M_k^t(o, a, \tilde{a}) - \left(\mathbb{O}_2 \mathbb{T}_1(\tilde{a}) \text{diag}(\mu_1) \mathbb{O}_1^\top\right) \right\|_2 &\leq \mathcal{O}\left(\sqrt{\frac{\ell}{k}}\right), \\
\left\| \frac{1}{k} N_k^h(a, \tilde{a}) - \mathbb{O}_h \mathbb{T}_{h-1}(\tilde{a}) \text{diag}(\mu_{h-1}^k) \mathbb{O}_{h-1}^\top \right\|_F &\leq \mathcal{O}\left(\sqrt{\frac{\ell}{k}}\right), \\
\left\| \frac{1}{k} M_k^h(o, a, \tilde{a}) - \mathbb{O}_{h+1} \mathbb{T}_h(a) \text{diag}(\mu_{h-1}^k) \mathbb{T}_{h-1}(\tilde{a}) \text{diag}(\mu_{h-1}^k) \mathbb{O}_{h-1}^\top \right\|_F &\leq \mathcal{O}\left(\sqrt{\frac{\ell}{k}}\right),
\end{align*}
\]
where
\[ \tau = \log(KAOH/\delta) \quad \text{and} \quad \mu_{h-1}^k = \frac{1}{k} \sum_{t=1}^{k} \mathbb{P}_{\pi_t}^k(s_{h-1} = \cdot) \quad 2 \leq h \leq H - 1. \]

Note that for all \( k \in [K], \mu_1^k = \mu_1 \) independent of \( \pi_1, \ldots, \pi_k \).

Now, with Corollary 13 we can prove the true parameter \( \theta^* \) always lies in the confidence sets for \( k \in [K] \) with high probability.

**Lemma 14.** Denote by \( \theta^* = (T, O, \mu_1) \) the the ground truth parameters of the POMDP. With probability at least \( 1 - \delta \), we have \( \theta^* \in \Theta_k \) for all \( k \in [K] \).

**Proof.** By the definition of \( b_0(\theta^*) \) and \( B_h(a, o; \theta^*) \), we have
\[
\begin{aligned}
b_0(\theta^*) &= \mathbb{O}_1 \mu_1, \\
\theta^* &= \left( \mathbb{O}_2 T_1(\bar{a}) \text{diag}(\mu_1) \mathbb{O}_1^\top \right)_o = B_1(\bar{a}, o; \theta^*) \mathbb{O}_1 \mu_1, \\
W &= B_h(a, o; \theta^*) \cdot V, \quad h \geq 2.
\end{aligned}
\]

where \( W \) and \( V \) are shorthands defined in Corollary 13.

It is easy to see \((*)\) and Corollary 13 directly imply \( \| n^k - b_0(\theta^*) \|_2 \leq O \left( \sqrt{k\tau} \right) \) and thus \( \theta^* \) satisfies Type-0 constraint. For other constraints, we have

**Type-I constraint:**
\[
\begin{aligned}
\| M^k_1(a, a, \bar{a}) - B_1(\bar{a}, o; \theta^*) N^k_1(a, \bar{a}) \|_2 &\leq \| M^k_1(a, a, \bar{a}) - k \left( \mathbb{O}_2 T_1(\bar{a}) \text{diag}(\mu_1) \mathbb{O}_1^\top \right)_o B_1(\bar{a}, o; \theta^*) \mathbb{O}_1 \mu_1 \|_2 \\
&\quad + k \| \left( \mathbb{O}_2 T_1(\bar{a}) \text{diag}(\mu_1) \mathbb{O}_1^\top \right)_o - B_1(\bar{a}, o; \theta^*) \mathbb{O}_1 \mu_1 \|_2 \\
&= \| M^k_1(a, a, \bar{a}) - k \left( \mathbb{O}_2 T_1(\bar{a}) \text{diag}(\mu_1) \mathbb{O}_1^\top \right)_o \|_2 + \| B_1(\bar{a}, o; \theta^*) \mathbb{O}_1 \mu_1 - N^k_1(a, \bar{a}) \|_2 \\
&\leq \| M^k_1(a, a, \bar{a}) - k \left( \mathbb{O}_2 T_1(\bar{a}) \text{diag}(\mu_1) \mathbb{O}_1^\top \right)_o \|_2 + \| B_1(\bar{a}, o; \theta^*) \mathbb{O}_1 \mu_1 \|_2 \| k \mathbb{O}_1 \mu_1 - N^k_1(a, \bar{a}) \|_2 \\
&\leq O \left( \frac{\sqrt{kS\tau}}{\alpha} \right)
\end{aligned}
\]

where the identity follows from \((*)\), and the last inequality follows from Corollary 13 and
\[
\| B_h(a, o; \theta^*) \|_2 = \| \mathbb{O}_{h+1} T_h(a) \text{diag}(\mathbb{O}_h(o|\cdot)) \mathbb{O}_h^\top \|_2 \\
\leq \frac{1}{\alpha} \| \mathbb{O}_{h+1} T_h(a) \text{diag}(\mathbb{O}_h(o|\cdot)) \|_2 \\
\leq \frac{\sqrt{\mathcal{S}}}{\alpha} \| \mathbb{O}_{h+1} T_h(a) \text{diag}(\mathbb{O}_h(o|\cdot)) \|_{1 \to 1} \leq \frac{\sqrt{\mathcal{S}}}{\alpha}.
\]

**Type-II constraint:** similarly, for \( h \geq 2 \), we have
\[
\begin{aligned}
\| B_h(a, o; \theta^*) N^k_h(a, \bar{a}) - M^k_h(a, o, \bar{a}) \|_F &\leq k \| B_h(a, o; \theta^*) \cdot V - W \|_F + \| B_h(a, o; \theta^*) (N^k_h(a, \bar{a}) - k V) \|_F + \| kW - M^k_h(a, o, \bar{a}) \|_F \\
&= \| B_h(a, o; \theta^*) (N^k_h(a, \bar{a}) - k V) \|_F + \| kW - M^k_h(a, o, \bar{a}) \|_F \\
&\leq \| B_h(a, o; \theta^*) \|_2 \| N^k_h(a, \bar{a}) - k V \|_F + \| kW - M^k_h(a, o, \bar{a}) \|_F \\
&\leq O \left( \frac{\sqrt{kS\tau}}{\alpha} \right).
\end{aligned}
\]
Therefore, we conclude that $\theta^* \in \Theta_k$ for all $k \in [K]$ with probability at least $1 - \delta$. □

Furthermore, with Corollary [13] we can prove the following bound for operator error.

**Lemma 15.** With probability at least $1 - \delta$, for all $k \in [K]$, $\hat{\theta} = (\hat{\Omega}, \hat{T}, \hat{\mu}_1) \in \Theta_{k+1}$ and $(o, a, \tilde{a}, h) \in \Theta \times \mathcal{A}^2 \times \{2, \ldots, H - 1\}$, we have

$$
\begin{align*}
\left\| b_0(\theta^*) - b_0(\hat{\theta}) \right\|_2 &\leq O\left( \sqrt{\frac{k}{k}} \right), \\
\left\| \left( B_1(\tilde{a}, o; \hat{\theta}) - B_1(\tilde{a}, o; \theta^*) \right) b_0(\theta^*) \right\|_2 &\leq O\left( \sqrt{\frac{S_k}{k\alpha^2}} \right)
\end{align*}
$$

$$
\sum_{s=1}^{S} \left\| \left( B_h(a, o; \hat{\theta}) - B_h(a, o; \theta^*) \right) \left( \hat{\Omega}_h \hat{T}_{h-1}(\tilde{a}) \right)_s \right\|_2 \sum_{t=1}^{k} \mathbb{P}_{\hat{\theta}}(\tilde{s}_{h-1} = s) \leq O\left( \sqrt{\frac{kS^2O_k}{\alpha^4}} \right),
$$

where $k = \log(KAOH/\delta)$.

**Proof.** For readability, we copy the following set of identities from Lemma [14] here,

$$
\begin{align*}
\mathbf{b}_0(\theta^*) &= \mathbf{O}_1\mu_1, \\
\left( \mathbf{O}_2\mathbf{T}_1(\tilde{a})\mathbf{diag}(\mu_1)\mathbf{O}_1^\top \right)_o &= \mathbf{B}_1(\tilde{a}, o; \theta^*)\mathbf{O}_1\mu_1, \\
\mathbf{W} &= \mathbf{B}_h(a, o; \theta^*) \cdot \mathbf{V}, \quad h \geq 2,
\end{align*}
$$

**Type-0 closeness:**

$$
\left\| b_0(\theta^*) - b_0(\hat{\theta}) \right\|_2 \leq \left\| \frac{1}{k} \mathbf{n}^k - b_0(\theta^*) \right\|_2 + \left\| \frac{1}{k} \mathbf{n}^k - b_0(\hat{\theta}) \right\|_2 \leq O\left( \sqrt{\frac{k}{k}} \right),
$$

where the last inequality follows from (*) and Corollary [14] and $\hat{\theta} \in \Theta_{k+1}$.

**Type-I closeness:** similarly, we have

$$
\begin{align*}
\left\| \left( B_1(\tilde{a}, o; \hat{\theta}) - B_1(\tilde{a}, o; \theta^*) \right) b_0(\theta^*) \right\|_2 &\leq \left\| \left( \mathbf{O}_2\mathbf{T}_1(\tilde{a})\mathbf{diag}(\mu_1)\mathbf{O}_1^\top \right)_o - \mathbf{B}_1(\tilde{a}, o; \theta^*)b_0(\theta^*) \right\|_2 \\
&\quad + \left\| \left( \mathbf{O}_2\mathbf{T}_1(\tilde{a})\mathbf{diag}(\mu_1)\mathbf{O}_1^\top \right)_o - \mathbf{B}_1(\tilde{a}, o; \hat{\theta})b_0(\theta^*) \right\|_2 \\
&= \left\| \left( \mathbf{O}_2\mathbf{T}_1(\tilde{a})\mathbf{diag}(\mu_1)\mathbf{O}_1^\top \right)_o - \mathbf{B}_1(\tilde{a}, o; \hat{\theta})b_0(\theta^*) \right\|_2 \\
&\leq \left\| \left( \mathbf{O}_2\mathbf{T}_1(\tilde{a})\mathbf{diag}(\mu_1)\mathbf{O}_1^\top \right)_o - \frac{1}{k} \mathbf{M}_1^k(o, a, \tilde{a}) \right\|_2 + \frac{1}{k} \left\| \mathbf{M}_1^k(o, a, \tilde{a}) - \mathbf{B}_1(\tilde{a}, o; \hat{\theta})\mathbf{N}_1^k(a, \tilde{a}) \right\|_2 \\
&\quad + \left\| \mathbf{B}_1(\tilde{a}, o; \hat{\theta}) \left( \frac{1}{k} \mathbf{N}_1^k(a, \tilde{a}) - b_0(\theta^*) \right) \right\|_2 \\
&\leq \left\| \left( \mathbf{O}_2\mathbf{T}_1(\tilde{a})\mathbf{diag}(\mu_1)\mathbf{O}_1^\top \right)_o - \frac{1}{k} \mathbf{M}_1^k(o, a, \tilde{a}) \right\|_2 + \frac{1}{k} \left\| \mathbf{M}_1^k(o, a, \tilde{a}) - \mathbf{B}_1(\tilde{a}, o; \hat{\theta})\mathbf{N}_1^k(a, \tilde{a}) \right\|_2 \\
&\quad + \left\| \mathbf{B}_1(\tilde{a}, o; \hat{\theta}) \right\|_2 \left\| \frac{1}{k} \mathbf{N}_1^k(a, \tilde{a}) - \mathbf{O}_1\mu_1 \right\|_2 \\
&\leq O\left( \sqrt{\frac{S_k}{k\alpha^2}} \right),
\end{align*}
$$

where the identity follows from (*) and the last inequality follows from Corollary [13] and $\hat{\theta} \in \Theta_{k+1}$.  

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Moreover, there exists some hard case where we have an almost matching lower bound

\[ \| \left( B_h(a, o; \hat{\theta}) - B_h(a, o; \theta^*) \right) V \|_F \]

\[ \leq \| W - B_h(a, o; \theta^*) V \|_F + \frac{1}{k} \| \mathcal{M}_h^k(o, a, \hat{a}) - \mathcal{M}_h^k(o, a, \hat{a}) \|_F \\
\quad + \frac{1}{k} \| B_h(a, o; \hat{\theta}) \mathbb{N}_h^k(a, \alpha) - \mathcal{M}_h^k(o, a, \alpha) \|_F + \| B_h(a, o; \hat{\theta}) \left( V - \frac{1}{k} \mathbb{N}_h^k(a, \alpha) \right) \|_F \\
\quad = \frac{1}{k} \| \mathcal{M}_h^k(o, a, \alpha) - \mathcal{M}_h^k(o, a, \alpha) \|_F + \frac{1}{k} \| B_h(a, o; \hat{\theta}) \mathbb{N}_h^k(a, \alpha) - \mathcal{M}_h^k(o, a, \alpha) \|_F \\
\quad + \| B_h(a, o; \hat{\theta}) \left( V - \frac{1}{k} \mathbb{N}_h^k(a, \alpha) \right) \|_F \\
\quad \leq O \left( \sqrt{\frac{S^0}{k\alpha^2}} \right) \tag{14} \]

where the identity follows from (\ast) and the last inequality follows from Corollary 13 and \( \hat{\theta} \in \Theta_{k+1} \).

Recall \( V = \Omega_h \mathcal{T}_{h-1}(\alpha) \text{diag}(\mu_{h-1}^k) \Omega_{h-1}^\top \) and utilize Assumption 1,

\[ \| \left( B_h(a, o; \hat{\theta}) - B_h(a, o; \theta^*) \right) V \|_F \geq \alpha \| \left( B_h(a, o; \hat{\theta}) - B_h(a, o; \theta^*) \right) \Omega_h \mathcal{T}_{h-1}(\alpha) \text{diag}(\mu_{h-1}^k) \|_F \]

\[ \geq \frac{\alpha}{\sqrt{SO}} \| \left( B_h(a, o; \hat{\theta}) - B_h(a, o; \theta^*) \right) \Omega_h \mathcal{T}_{h-1}(\alpha) \text{diag}(\mu_{h-1}^k) \|_1 \]

\[ = \frac{\alpha}{k\sqrt{SO}} \sum_{s=1}^S \| \left( B_h(a, o; \hat{\theta}) - B_h(a, o; \theta^*) \right) \Omega_h \mathcal{T}_{h-1}(\alpha) \|_1 \sum_{t=1}^k \| P_{\theta^*}^{(s)}(s_{h-1} = s) \|. \]

Plugging it back into (14) completes the whole proof. \( \square \)

### C.4 Proof of Theorem 3

In order to utilize Lemma 15 to bound the operator error in Lemma 10, we need the following algebraic transformation. Its proof is postponed to Appendix E.

**Lemma 16.** Let \( z_k \in [0, C_z] \) and \( w_k \in [0, C_w] \) for \( k \in \mathbb{N} \). Define \( S_k = \sum_{j=1}^k w_j \) and \( S_0 = 0 \). If \( z_k S_{k-1} \leq C_z C_w C_0 \sqrt{k} \) for any \( k \in [K] \), we have

\[ \sum_{k=1}^K z_k w_k \leq 2C_z C_w (C_0 + 1) \sqrt{K} \log(K). \]

Moreover, there exists some hard case where we have a almost matching lower bound \( O \left( C_z C_w C_0 \sqrt{K} \right) \).

Now, we are ready to prove the main theorem based on Lemma 10, Lemma 15, and Lemma 16.

**Theorem 3.** For any \( \varepsilon \in (0, H] \), there exists \( K_{\max} = \text{poly}(H, S, A, O, \alpha^{-1})/\varepsilon^2 \) and an absolute constant \( c_1 \), such that for any POMDP that satisfies Assumption 7 if we set hyperparameters \( \beta_h = c_1 \sqrt{\log(KAOH)} \), \( \gamma_h = \sqrt{S \beta_h / \alpha} \), and \( K \geq K_{\max} \), then the output policy \( \hat{\pi} \) of Algorithm 7 will be \( \varepsilon \)-optimal with probability at least 2/3.
We can bound the first two terms by Lemma 15, and obtain that with probability at least

where the identity follows from Fact 18.

By Lemma 14, we have \( \theta^* \in \Theta_k \) for all \( k \in [K] \) with probability at least \( 1 - \delta \). Recall that \( (\pi_k, \theta_k) = \arg\max_{\pi, \theta \in \Theta_k} V^\pi(\theta) \), so with probability at least \( 1 - \delta \), we have

\[
\sum_{k=1}^K \left( V^\pi^*(\theta^*) - V^{\pi_k}(\theta^*) \right) \\
\leq \sum_{k=1}^K \left( V^{\pi_k}(\theta_k) - V^{\pi_k}(\theta^*) \right) \\
\leq H \sum_{k=1}^K \sum_{\tau_{H-1} \in \Gamma(\pi_k, H-1)} \left\| \mathbb{P}^k_\theta ([o_H, \tau_{H-1}]) - \mathbb{P}^k_\theta([o_H, \tau_{H-1}]) \right\|_1 \\
= H \sum_{k=1}^K \sum_{\tau_{H-1} \in \Gamma(\pi_k, H-1)} \left\| b(\tau_{H-1}; \theta^*) - b(\tau_{H-1}; \theta_k) \right\|_1,
\]

(15)

where the identity follows from Fact 18.

Applying Lemma 10, we have

\[
\sum_{\tau_{H-1} \in \Gamma(\pi_k, H-1)} \left\| b(\tau_{H-1}; \theta^*) - b(\tau_{H-1}; \theta_k) \right\|_1 \\
\leq \sqrt{S} \left\| b_0(\theta^*) - b_0(\theta_k) \right\|_1 + \sqrt{S} \sum_{(a, o) \in \mathcal{A} \times \mathcal{O}} \left\| (B_1(a, o; \theta_k) - B_1(a, o; \theta^*)) b_0(\theta^*) \right\|_1 \\
+ \sqrt{S} \sum_{h=2}^{H-1} \sum_{(a, \bar{a}, o) \in \mathcal{A}^2 \times \mathcal{O}} \sum_{s=1}^S \left\| (B_h(a, o; \theta_k) - B_h(a, o; \theta^*)) (\bigcirc_h \mathbb{T}_{h-1}(\bar{a})) s \right\|_1 \mathbb{P}^k_\theta(s_{h-1} = s).
\]

(16)

We can bound the first two terms by Lemma 13, and obtain that with probability at least \( 1 - \delta \),

\[
H \sum_{k=1}^K J_k \leq C_{\text{poly}} \sqrt{K} \log(K).
\]

(17)

Plugging (16) and (17) into (15), we obtain

\[
\sum_{k=1}^K \left( V^\pi^*(\theta^*) - V^{\pi_k}(\theta^*) \right) \leq C_{\text{poly}} \sqrt{K} \log(K) + \\
C_{\text{poly}} \max_{s, o, a, h} \sum_{k=1}^K \left\| (B_h(a, o; \theta_k) - B_h(a, o; \theta^*)) (\bigcirc_h \mathbb{T}_{h-1}(\bar{a})) s \right\|_1 \mathbb{P}^k_\theta(s_{h-1} = s).
\]

(18)

It remains to bound the second term.

By Lemma 15, with probability at least \( 1 - \delta \), for all \( k \in [K] \), \( \theta_k \in \Theta_k \) and \( (s, o, a, \bar{a}, h) \in \mathcal{I} \times \mathcal{O} \times \mathcal{A}^2 \times \{2, \ldots, H - 1\} \), we have

\[
\left\| (B_h(a, o; \theta_k) - B_h(a, o; \theta^*)) (\bigcirc_h \mathbb{T}_{h-1}(\bar{a})) s \right\|_1 \sum_{\tilde{w}_t} \mathbb{P}^k_\theta(s_{h-1} = s) \leq C_{\text{poly}} \sqrt{K} \log(K).
\]

(19)
Invoking Lemma 16 with (19), we obtain

$$\sum_{k=1}^{K} w_k z_k \leq C_{\text{poly}} \sqrt{K} \log^2(K). \quad (20)$$

Plugging (20) back into (18), choosing $\delta = 1/10$ and outputting a policy from $\{\pi_1, \ldots, \pi_K\}$ uniformly at random complete the proof. \hfill \Box

## D Learning POMDPs with Deterministic Transition

In this section, we introduce a computationally and statistically efficient algorithm for POMDPs with deterministic transition. A sketched proof is provided.

### Algorithm 2 Learning POMDPs with Deterministic Transition

1: initialize $N = C \log(HSA/p)/(\min\{\varepsilon/(\sqrt{O}H), \xi\})^2$, $n_h = 1(h = 1)$ for all $h \in [H]$.
2: for $h = 1, \ldots, H - 1$ do
3: for $(s, a) \in [n_h] \times A$ do
4: Reset $z \leftarrow 0_{O \times 1}$ and $t \leftarrow n_{h+1} + 1$
5: for $i \in [N]$ do
6: execute policy $\pi_h(s)$ from step 1 to step $h - 1$, take action $a$ at $h^{th}$ step and observe $o_{h+1}$
7: $z \leftarrow z + \frac{1}{\sqrt{C}} e_{o_{h+1}}$
8: for $s' \in [n_{h+1}]$ do
9: if $\|\phi_{h+1,s'} - z\|_2 \leq 0.5 \xi$ then
10: $t \leftarrow s'$
11: if $t = n_{h+1} + 1$ then
12: $n_{h+1} \leftarrow n_{h+1} + 1$
13: $\phi_{h+1,n_{h+1}} \leftarrow z$ and $\pi_{h+1}(n_{h+1}) \leftarrow a \circ \pi_h(s)$
14: Set the $s^{th}$ column of $\bar{T}_{h,a}$ to be $e_t$
15: output $\hat{\mu}_0 = e_1$ and $\{n_h, \{\bar{T}_{h,a}\}_{a \in A} \text{ and } \{\phi_{h,i}\}_{i \in [n_h]} : h \in [H]\}$

### Theorem 4. For any $p \in (0, 1]$, there exists an algorithm such that for any POMDP with deterministic transitions that satisfies Assumption 2 within $O \left( H^2SA \log(HSA/p)/(\min\{\varepsilon/(\sqrt{O}H), \xi\})^2 \right)$ samples and computations, the output policy of the algorithm is $\varepsilon$-optimal with probability at least $1 - p$.

**Proof.** The algorithm works by inductively finding all the states we can reach at each step, utilizing the property of deterministic transition and good separation between different observation vectors. We sketch a proof based on induction below.

We say a state $s$ is $h$-step reachable if there exists a policy $\pi$ s.t. $\mathbb{P}^\pi(s_h = s) = 1$. In our algorithm, we use $n_h$ to denote the number of $h$-step reachable states. All the policies mentioned below is a sequence of fixed actions (independent of observations).

Suppose at step $h$, there are $n_h$ $h$-step reachable states and we can reach the $s^{th}$ one of them at the $h^{th}$ step by executing a known policy $\pi_h(s)$. Note that for every state $s'$ that is $(h + 1)$-step reachable, there must exist some state $s$ and action $a$ s.t. $s$ is $h$-step reachable and $T_h(s' | s, a) = 1$. Therefore, based on our induction assumption, we can reach all the $(h + 1)$-step reachable states by executing all $a \circ \pi_h(s)$ for $(a, s) \in A \times [n_h]$.

Now the problem is how to tell if we reach the same state by executing two different $a \circ \pi_h(s)$'s. The solution is to look at the distribution of $o_{h+1}$. Because the POMDP has deterministic transition, we always reach the
same state when executing the same \( a \circ \pi_h(s) \) and hence the distribution of \( o_{h+1} \) is exactly the distribution of observation corresponding to that state. By Hoeffding’s inequality, for each fixed \( a \circ \pi_h(s) \), we can estimate the distribution of \( o_{h+1} \) with \( \ell_2 \)-error smaller than \( \xi/8 \) with high probability using \( N \geq \tilde{\Omega}(1/\xi^2) \) samples. Since the observation distributions of two different states have \( \ell_2 \)-separation no smaller than \( \xi \), we can tell if two different \( a \circ \pi_h(s) \)'s reach the same state by looking at the distance between their distributions of \( o_{h+1} \). For those policies reaching the same state, we only need to keep one of them, so there are at most \( S \) policies kept \( (n_{h+1} \leq S) \).

By repeating the argument inductively from \( h = 1 \) to \( h = H \), we can recover the exact transition dynamics \( T_h(\cdot \mid s, a) \) and get an high-accuracy estimate of \( \Omega_h(\cdot \mid s) \) for every \( h \)-step reachable state \( s \) and all \( (h, a) \in [H] \times \mathcal{A} \) up to permutation of states. Since the POMDP has deterministic transition, we can easily find the optimal policy of the estimated model by dynamic programming.

The \( \epsilon \)-optimality simply follows from the fact that when \( N \geq \tilde{\Omega}(H^2O/\epsilon^2) \), we have the estimated distribution of observation for each state being \( \mathcal{O}(\epsilon/H) \) accurate in \( \ell_1 \)-distance for all reachable states. This implies that the optimal policy of the estimated model is at most \( \mathcal{O}(\epsilon/H) \times H = \mathcal{O}(\epsilon) \) suboptimal. The overall sample complexity follows from our requirement \( N \geq \max\{\tilde{\Omega}(H^2O/\epsilon^2), \tilde{\Omega}(1/\xi^2)\} \), and the fact we need to run \( N \) episodes for each \( h \in [H], s \in \mathcal{S}, a \in \mathcal{A} \).

## E Auxiliary Results

### E.1 Basic facts about POMDPs and the operators

In this section, we provide some useful facts about POMDPs. Since their proofs are quite straightforward, we omit them here.

The following fact gives two linear equations the operators always satisfy. Its proof simply follows from the definition of the operators and Fact(11)

**Fact 17.** In the same setting as Fact(11) suppose Assumption(11) holds, then we have

\[
\begin{align*}
\mathbb{P}(o_h = \cdot, o_{h-1} = \cdot) & = B_h(\vec{a}, o; \theta)\mathbb{P}(o_{h-1} = \cdot), \\
\mathbb{P}(o_{h+1} = \cdot, o_h = a, o_{h-1} = \cdot) & = B_h(a, o; \theta)\mathbb{P}(o_h = \cdot, o_{h-1} = \cdot).
\end{align*}
\]

The following fact relates (unnormalized) belief states to distributions of observable sequences. Its proof follows from simple computation using conditional probability formula and \( \Omega^\top_h \Omega_h = I_S \).

**Fact 18.** For any POMDP(\( \theta \)) satisfying Assumption(7) deterministic policy \( \pi \) and \( [o_h, \tau_{h-1}] \in \mathcal{S} \times \Gamma(\pi, h - 1) \), we have

\[
e_{o_h}^\top b(\tau_{h-1}; \theta) = \mathbb{P}_{\theta}^\pi([o_h, \tau_{h-1}]),
\]

where \( \mathbb{P}_{\theta}^\pi([o_h, \tau_{h-1}]) \) is the probability of observing \( [o_h, \tau_{h-1}] \) when executing policy \( \pi \) in POMDP(\( \theta \)).

### E.2 Proof of Lemma(16)

**Proof.** WLOG, assume \( C_z = C_w = 1 \). Let \( n = \min\{k \in [K] : S_k \geq 1\} \). We have

\[
\sum_{k=1}^{K} z_k w_k = \sum_{k=1}^{n} z_k w_k + \sum_{k=n+1}^{K} z_k w_k \leq \sum_{k=1}^{n} w_k + \sum_{k=n+1}^{K} z_k w_k
\]

\[
= S_n + \sum_{k=n+1}^{K} z_k w_k \leq 2 + \sum_{k=n+1}^{K} z_k w_k.
\]
It remains to bound the second term. Using the condition that $z_k S_{k-1} \leq C_0 \sqrt{k}$ for all $k \in [K]$, we have $z_k \leq \frac{C_0 \sqrt{K}}{S_{k-1}}$ for all $k \in [K]$ and $i \in [m]$. Therefore,

\[
\sum_{k=n+1}^{K} z_k w_k \leq \sum_{k=n+1}^{K} C_0 \sqrt{k} \frac{w_k}{S_{k-1}} \leq C_0 \sqrt{K} \sum_{k=n+1}^{K} \frac{w_k}{S_{k-1}} \leq (a) 2C_0 \sqrt{K} \sum_{k=n+1}^{K} \log \left( \frac{S_k}{S_{k-1}} \right) = 2C_0 \sqrt{K} \log \left( \frac{S_K}{S_n} \right) \leq 2C_0 \sqrt{K} \log(K),
\]

where (a) follows from $x \leq 2 \log(x + 1)$ for $x \in [0, 1]$. 

\[\Box\]