Research Article

Stability of the Generalized Polar Decomposition Method for the Approximation of the Matrix Exponential

Elham Nobari and S. Mohammad Hosseini

Department of Mathematics, Tarbiat Modares University, P.O. Box 14115-175, Tehran, Iran
Address correspondence to S. Mohammad Hosseini, hossei_m@modares.ac.ir

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Abstract Generalized polar decomposition method (or briefly GPD method) has been introduced by Munthe-Kaas and Zanna [5] to approximate the matrix exponential. In this paper, we investigate the numerical stability of that method with respect to roundoff propagation. The numerical GPD method includes two parts: splitting of a matrix $Z \in g$, a Lie algebra of matrices and computing $\exp(Z)v$ for a vector $v$. We show that the former is stable provided that $\|Z\|$ is not so large, while the latter is not stable in general except with some restrictions on the entries of the matrix $Z$ and the vector $v$.

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1 Introduction

Since exponential matrix naturally appears in the analytical solutions of many differential equations, it is important to present the numerical approximate methods for computing exponential matrices in applied problems. In many cases that the different equations are modeled on a matrix Lie group, we need that the approximated matrix to remain in the same Lie group. Most of the well-known standard methods [3] do not satisfy this condition except in some special cases, like $3 \times 3$ antisymmetric matrices [2].

Recently, Munthe-Kaas and Zanna have introduced a geometric numerical method, called generalized polar decomposition method, or briefly GPD method [5]. The main advantage of this method among others is that it preserves the approximated object in the mentioned Lie group. In the GPD method, the Lie group $G$ and its Lie algebra $g$ are splitted to simpler ones, such that the computation of $\exp Z$ will be easier in those subgroups and subalgebras. There are two numerical algorithms in GPD method. The first one (Algorithm 1) splits the matrix $Z$ and in the second one (Algorithm 2), $\exp(Z)v$ is computed, where $v$ is an arbitrary vector.

As far as we know, no proof of the stability of the GPD method, with respect to propagation of roundoff error, has appeared in the literature. In this paper, we study the stability of the method for the so-called polar-type splitting order-two algorithm that has been introduced in [5]. We show that Algorithm 1 which is a simpler algorithm is stable, provided the norm $\|Z\|$ of the matrix $Z$ and in the second one (Algorithm 2), $\exp(Z)v$ is computed, where $v$ is an arbitrary vector.

We introduce a sufficient condition for this aim. More precisely, given an $n \times n$ matrix $Z$, let

$$a_j = [Z_{j+1,j}, Z_{j+2,j}, \ldots, Z_{n,j}], \quad b_j = [Z_{j,j+1}, Z_{j,j+2}, \ldots, Z_{j,n}]^T$$

for $j = 1, 2, \ldots, n-1$ and let

$$k_0 = \min \{ \|b_j^T a_j\| : b_j^T a_j \neq 0, \ 1 \leq j \leq n-1 \},$$

then Algorithm 2 is stable under the following conditions:

$$\frac{\|Z\|}{k_0} < \left( \frac{1 - \alpha}{1 + \alpha} \right) \left( \frac{1 - n\alpha}{n^2\alpha} \right), \quad (1.1b)$$

where $\alpha$ is the unit roundoff error.

The content of the paper is as follows. In Section 2, some preliminaries concerning the general concept of error analysis and GPD method and also details of the above-mentioned algorithms are given. In Section 3, the error analysis of the Algorithm 1 and in Section 4, the error analysis of Algorithm 2 are addressed. Section 5 is devoted to sensibility analysis of Algorithm 2.
2 Preliminaries

2.1 Error analysis

In this section, we recall some standard definitions and lemmas, which are applied in the other sections. Elementary operators like $+, \cdot, \times,$ and $/$ as well as $\sqrt{\cdot}, \sin,$ and $\cos$ admit the following floating point arithmetic:

$$\text{fl}(x \text{ op } y) = (x \text{ op } y)(1 + \delta), \quad |\delta| \leq u,$$

where $\text{op}$ is an elementary operation, and $u$ is the unit roundoff error.

**Lemma 1** (see [1]). If $|\delta_i| \leq u$ and $\rho_i = \pm 1$ for $1 \leq i \leq n$ and $nu < 1$, then

$$\prod_{i=1}^{n} (1 + \delta_i)^{\rho_i} = 1 + \theta_n,$$

such that

$$|\theta_n| \leq \frac{nu}{1 - nu} := \gamma_n.$$

The following technical properties are satisfied by the sequence $\{\gamma_n\}$:

$$\gamma_j + \gamma_k + \gamma_j \gamma_k \leq \gamma_{j+k}, \quad (j, k \geq 1), \quad (2.1)$$

$$c \gamma_n \leq \gamma_{cn}, \quad (c \geq 1, n > 1). \quad (2.2)$$

Denoting the calculated value of any variable $x$ by $\hat{x}$ and $|A| = \begin{bmatrix} |a_{ij}| \end{bmatrix}$ for any matrix $A = [a_{ij}]$, we have the following lemma.

**Lemma 2** (see [1,4]). If $y = (c - \sum_{i=1}^{k-1} a_i b_i)/b_k$ is evaluated in floating point arithmetic, then, no matter what is the order of evaluation, there exist constants $\theta_k^{(i)}$, $i = 0, 1, \ldots, k - 1$, such that

$$b_k \hat{y}(1 + \theta_k^{(0)}) = c - \sum_{i=1}^{k-1} a_i b_i (1 + \theta_k^{(i)}), \quad (2.3)$$

where $|\theta_k^{(i)}| \leq \gamma_k$ for all $i$. If $b_k = 1$, so that there is no division, then $|\theta_k^{(i)}| \leq \gamma_{k-1}$ for all $i$.

Also for matrix multiplication, we have

$$\text{fl}(AB) = AB + \Delta, \quad |\Delta| \leq \gamma_n |A||B|, \quad A \in \mathbb{R}^{m \times n}, B \in \mathbb{R}^{n \times p}. \quad (2.4)$$

For any matrix $A = [a_{ij}]$, let $\|A\| = \max_{i,j} |a_{ij}|$. By this norm, we have

$$\|AB\| \leq n\|A\||\|B\| \quad (2.5)$$

for any $A \in \mathbb{R}^{m \times n}$ and $B \in \mathbb{R}^{n \times p}$.

For any two vectors $a, b \in \mathbb{R}^n$, let

$$a \ast b = (a_1 b_1, \ldots, a_n b_n) \in \mathbb{R}^n. \quad (2.6)$$

So, using this notation, we can write

$$\text{fl}(\lambda a) = \lambda a \ast (1 + u), \quad |u| \leq u1, \quad 1 = (1, \ldots, 1) \in \mathbb{R}^n, \lambda \in \mathbb{R}. \quad (2.7)$$

Similar to Lemma 1, one can prove the following lemma.

**Lemma 3.** If $|u_i| \leq u1$, for $1 \leq i \leq n$ and $nu < 1$, then

$$(1 + u_1) \ast \cdots \ast (1 + u_n) = 1 + h_n, \quad (2.8)$$

such that

$$|h_n| \leq \gamma_n 1.$$
2.2 GPD method

Suppose $G$ is a finite dimensional Lie group and $\mathfrak{g}$ is its Lie algebra and let $\sigma : G \rightarrow G$, $\sigma \neq id$ be an involutive automorphism, that is, one-to-one smooth map such that:

$$\sigma(x \cdot y) = \sigma(x) \cdot \sigma(y), \quad \forall x, y \in G,$$

$$\sigma(\sigma(x)) = x, \quad \forall x \in G.$$  

It is well known that for sufficiently small $t$ and for any $z = \exp(tZ) \in G$, where $Z \in \mathfrak{g}$, we can write

$$z = xy,$$

(2.9)

where $\sigma(x) = x^{-1}$ and $\sigma(y) = y$. The decomposition (2.9) is called the Generalized Polar Decomposition (GPD) of $z$. On the Lie algebra $\mathfrak{g}$, the involutive automorphism is induced by $\sigma$:

$$d\sigma(Z) = \frac{d}{dt} \bigg|_{t=0} \sigma(\exp tZ),$$

hence $d\sigma$ defines a splitting of $\mathfrak{g}$ into the direct sum of two linear spaces:

$$\mathfrak{g} = \mathfrak{p} \oplus \mathfrak{k},$$

where $\mathfrak{k} = \{ Z \in \mathfrak{g} : d\sigma(Z) = Z \}$ and $\mathfrak{p} = \{ Z \in \mathfrak{g} : d\sigma(Z) = -Z \}$. In fact, $\mathfrak{k}$ is a Lie subalgebra of $\mathfrak{g}$, but $\mathfrak{p}$ only admits Lie triple system property:

$$A, B, C \in \mathfrak{p} \implies [A, [B, C]] \in \mathfrak{p},$$

where $[A, B]$ is the standard commutator on Lie algebra $\mathfrak{g}$. Now, let

$$P = \Pi_\mathfrak{p}(Z), \quad K = \Pi_\mathfrak{k}(Z),$$

where the maps $\Pi_\mathfrak{p} : \mathfrak{g} \rightarrow \mathfrak{p}$ and $\Pi_\mathfrak{k} : \mathfrak{g} \rightarrow \mathfrak{k}$ are the canonical projection maps. It can be easily verified that every $Z \in \mathfrak{g}$ can be uniquely decomposed into $Z = P + K$, where

$$P = \Pi_\mathfrak{p}(Z) = \frac{1}{2}(Z - d\sigma(Z)), \quad K = \Pi_\mathfrak{k}(Z) = \frac{1}{2}(Z + d\sigma(Z)),$$

(see [5]). Also, $\mathfrak{k}$ and $\mathfrak{p}$ satisfy

$$[\mathfrak{k}, \mathfrak{k}] \subseteq \mathfrak{k}, \quad [\mathfrak{k}, \mathfrak{p}] \subseteq \mathfrak{k}, \quad [\mathfrak{p}, \mathfrak{p}] \subseteq \mathfrak{k}.$$ 

If, in the decomposition (2.9), we consider $x = \exp X(t)$ and $y = \exp Y(t)$, where $X(t) \in \mathfrak{p}$ and $Y(t) \in \mathfrak{k}$, for sufficiently small $t$, then it can be written as follows:

$$X(t) = \sum_{i=1}^{\infty} X_i t^i, \quad Y(t) = \sum_{i=1}^{\infty} Y_i t^i,$$

where the coefficients $X_i$ and $Y_i$ can be found in [5]. The first terms in the expansions of $X(t)$ and $Y(t)$ are

$$X = Pt - \frac{1}{2} [P, K]^2 \frac{1}{6} [K, [P, K]] t^3 + \left( \frac{1}{24} [P, [P, [P, K]]] - \frac{1}{24} [K, [K, [P, K]]] \right) t^4$$

$$+ \left( \frac{7}{360} [K, [P, [P, K]]] - \frac{1}{120} [K, [K, [P, K]]] - \frac{1}{180} ([P, K], [P, [P, K]]) \right) t^5 + O(t^6),$$

(2.10)

$$Y = Kt - \frac{1}{12} [P, [P, K]] t^3$$

$$+ \left( \frac{1}{120} [P, [P, [P, K]]] + \frac{1}{720} [K, [K, [P, K]]] - \frac{1}{240} ([P, K], [K, [P, K]]) \right) t^5 + O(t^7).$$

(2.11)

Considering $\sigma_1, \sigma_2, \ldots, \sigma_m$ as elements of a sequence of involutive automorphisms on $G$, $\sigma_1$ induces a decomposition $\mathfrak{g} = \mathfrak{p}_1 \oplus \mathfrak{k}_1$ and approximation:

$$\exp(tZ) \approx \exp \left( X^{[1]}(t) \right) \exp \left( Y^{[1]}(t) \right).$$
where $X^{[1]}$ and $Y^{[1]}$ are suitable truncations of (2.10) and (2.11), respectively. Then $t_2$ can be decomposed by $\sigma_2$ and so we have $t_1 = p_2 \oplus t_2$ and

$$\exp \left( Y^{[1]}(t) \right) \approx \exp \left( X^{[2]}(t) \right) \exp \left( Y^{[2]}(t) \right).$$

Repeating this process $m$ times, $(m \geq 1)$ yields the following:

$$g = p_1 \oplus p_2 \oplus \cdots \oplus p_m \oplus t_m, \quad \exp(tZ) \approx \exp \left( X^{[1]}(t) \right) \cdots \exp \left( X^{[m]}(t) \right) \exp \left( Y^{[m]}(t) \right). \quad (2.12)$$

The number $m$ is chosen appropriately such that the right-hand side terms can be easily computed. If the expansions (2.10) are truncated at order two, then the two following algorithms based on the iterated generalized polar decomposition (2.12) will be obtained. The first algorithm splits the matrix $Z$, while the second one approximates $\exp(Z)\mathbf{v}$.

**Algorithm 1** (Polar-type splitting, order two).

% Purpose: 2nd order approximation of the splitting (2.12)
% In: $n \times n$ matrix $Z$
% Out: $Z$ overwritten with the nonzero elements of $X^{[1]}$ and $Y^{[m]}$ as:
% $Z(i+1:n,i) = X^{[i]}(i+1:n,i)$, $Z(i,i+1:n) = X^{[i]}(i,i+1:n)$
% $\text{diag}(Z) = \text{diag}(Y^{[m]})$

for $j = 1 : n - 1$
- $a_j := Z(j+1:n,j)$
- $b_j := Z(j,j+1:n)^T$
- $\overline{K}_j := Z(j+1:n,j+1:n)$
- $c_j := z_{ij}a_j - \overline{K}_j a_j$
- $d_j := -z_{ij}b_j + \overline{K}_j b_j$
- $Z(j+1:n,j) := a_j - \frac{1}{2}c_j$
- $Z(j,j+1:n) := (b_j - \frac{1}{2}d_j)^T$
end

**Algorithm 2** (Polar-type approximation).

% Purpose: Computing the approximant (2.12) applied to a vector $\mathbf{v}$
% In: $\mathbf{v}$: $n$-vector
% $Z$ : $n \times n$ matrix containing the nonzero elements of $X^{[i]}$ and $Y^{[m]}$ as:
% $Z(i+1:n,i) = X^{[i]}(i+1:n,i)$, $Z(i,i+1:n) = X^{[i]}(i,i+1:n)$
% $\text{diag}(Z) = \text{diag}(Y^{[m]})$
% Out: $\mathbf{v} := \exp(X^{[1]}) \cdots \exp(X^{[m]}) \exp(Y^{[m]})\mathbf{v}$

for $k = 1 : n$
- $v_k := \exp(z_{k,k})v_k$
end

for $j = n - 1 : -1$
- $a_j := [0; Z(j+1:n,j)]$
- $b_j := [0; Z(j,j+1:n)^T]$
- $v_{\text{old}} := \mathbf{v}(j:n)$
- $\alpha_j := \sqrt{b_j^T a_j}$, $\beta_j := \frac{\sinh \alpha_j}{\alpha_j}$, $\lambda_j := 2 \sinh^2 \left( \frac{\alpha_j}{2} \right)$
- $D := \begin{pmatrix} 0 & 1 \\ \alpha_j^2 & 0 \end{pmatrix}$
- $\varphi(D) := \lambda_j D^{-1} + \beta_j I$
- $w := \varphi(D) \begin{pmatrix} v_{\text{old}} \\ v_{\text{new}} \end{pmatrix}$
- $v_{\text{new}} := [a_j, e_1]w = w_1 a_j + w_2 e_1$
- $\mathbf{v}(j:n) = v_{\text{old}} + v_{\text{new}}$
end
3 Error analysis of Algorithm 1

In Algorithm 1, we denote by \( Z \) simultaneously the input data matrix and the splitted output matrix. In this algorithm, the calculated value of \( c_j \) (\( 1 \leq j \leq n-1 \)) is

\[
\hat{c}_j = (z_{j,j} a_j * (1 + u_j^1) - K_j a_j - u_j^2) * (1 + u_j^3),
\]

where \( |u_j^i| \leq u1 \) (\( i = 1, 3 \)) and \( |u_j^2| \leq \gamma_n \) \( |K_j a_j| \). From (2.8), we have \( (1 + u_j^1) * (1 + u_j^3) = 1 + h_2 \), where \( |h_2| \leq \gamma_2 \).

Hence,

\[
\hat{c}_j = c_j + u_j,
\]

where

\[
\begin{align*}
|u_j^1| & \leq \gamma_2 \sum z_{j,j} |a_j| + u |K_j a_j| + \gamma_n \sum |a_j| * (1 + u1).
\end{align*}
\]

(3.1)

By applying the norm of the matrix \( Z \) this reduces to

\[
|u_j| \leq (\gamma_2 + u + \gamma_n \sum u1) \|Z\|^2 1.
\]

Now, similarly for computing \( d_j \), we have

\[
\hat{d}_j = d_j + u_j,
\]

where

\[
|u_j^1| \leq (\gamma_2 + u + \gamma_n \sum u1) \|Z\|^2 1.
\]

(3.2)

Therefore, the entries of the matrix \( Z \) change as follows:

\[
\|a_j - \frac{1}{2} c_j\| = \|a_j - \frac{1}{2} (c_j + u_j) * (1 + u_j^3)\| * (1 + u_j^3) = \|a_j - \frac{1}{2} c_j\| + p_j,
\]

where \( |u_j^i| \leq u1 \) (\( i = 4, 5 \)) and

\[
p_j = a_j * u_j^3 - \frac{1}{2} c_j * h_2 - \frac{1}{2} u_j * (1 + h_2).
\]

So from (3.1),

\[
|p_j| \leq u |a_j| + \frac{1}{2} |c_j| \|z_{j,j} a_j| + (u + \gamma_n \sum u1) |K_j a_j|).
\]

If we replace the value of \( c_j \) from Algorithm 1 in the above relation, we will have

\[
|p_j| \leq u |a_j| + \left( \gamma_2 + \frac{1}{2} |z_{j,j} a_j| + \left( \frac{1}{2} \gamma_2 + \frac{1}{2} (u + \gamma_2) (u + \gamma_n \sum u1) \right) |K_j a_j|\right),
\]

and finally

\[
|p_j| \leq \left( u \|Z\| + \left( \gamma_2 + \frac{1}{2} \gamma_2 + \frac{1}{2} (u + \gamma_2) (u + \gamma_n \sum u1) \right) \|Z\|^2 \right) 1.
\]

Since

\[
\gamma_n = \frac{(n-j)u}{1-(n-j)u} \leq \frac{(n-j)u}{1-nu} = \frac{n-j}{n} \gamma_n
\]

for \( nu < 1 \) and from property (2.1), the upper bound of \( p_j \) will be

\[
|p_j| \leq \left( \|Z\| u + \left( \frac{2}{1-2u} + \frac{n-j+2}{2(1-nu)} + \frac{1}{2} \right) \|Z\|^2 u + O(u^2) \right) 1.
\]
The greatest value of the right-hand side occurs at $j = 1$, hence for all $j$,

$$|p_j| \leq \left( \|Z\|u + \left( \frac{2}{1 - 2u} + \frac{n + 1}{2(1 - nu)} + \frac{1}{2} \right)\|Z\|^2u + O(u^2) \right)1.$$  \hspace{1cm} (3.3)

Similarly, if

$$\text{fl} \left( b_j - \frac{1}{2}d_j \right) = \left( b_j - \frac{1}{2}d_j \right) + q_j,$$

then

$$|q_j| \leq \left( \|Z\|u + \left( \frac{2}{1 - 2u} + \frac{n + 1}{2(1 - nu)} + \frac{1}{2} \right)\|Z\|^2u + O(u^2) \right)1.$$ \hspace{1cm} (3.4)

However, since the diagonal entries do not change during this algorithm, (3.3) and (3.4) show the stability of the algorithm whenever $\|Z\|$ is not so large.

4 Error analysis of Algorithm 2

In the GPD method, the matrix $Z$ is splitted in Algorithm 1, and the $\exp(Z)v$ is calculated in Algorithm 2. We denote the splitted matrix of $Z$ with $Z = \left( \pi_{i,j} \right)$.

4.1 Error analysis of $\alpha_j$

Let $k_j = b^T_ja_j$ and $\alpha_j = \sqrt{|k_j|}$, here we have assumed $k_j > 0$. If $k_j$ is not a positive number, by considering $\alpha_j = \sqrt{|k_j|}$ when $k_j < 0$ and by letting $\frac{\sinh \alpha_j}{\alpha_j} = 1$ when $k_j = 0$, it is easily seen that of all the arguments and the upper bounds presented in the paper remain valid.

From Lemma 2, we have

$$\text{fl} \left( k_j \right) = \sum_{i=j+1}^{n} \pi_{j,i}\pi_{i,j} \left( 1 + \theta_{n-j}^{(i)} \right),$$

where $|\theta_{n-j}^{(i)}| \leq \gamma_{n-j}, \quad (1 \leq i \leq n)$. So we can write

$$\text{fl} \left( k_j \right) = k_j + \sum_{i=j+1}^{n} \pi_{j,i}\pi_{i,j}\theta_{n-j}^{(i)} = k_j + t,$$

where

$$|t| \leq (n - j)\gamma_{n-j}\|Z\|^2.$$ \hspace{1cm} (4.1)

Then by assuming $\|Z\| \leq 1$ and from condition (1.1), we have

$$|t| \leq n\gamma_n \leq \frac{1 - u}{1 + u}k_0 \leq k_0,$$

where $k_0 = \min \{ |b_j^T a_j|; b_j^T a_j \neq 0, 1 \leq j \leq n - 1 \}$. Hence $k_j + t$ is nonnegative and

$$\text{fl} \left( \alpha_j \right) = \sqrt{k_j + t} + (1 + \delta),$$

where $\delta$, with $|\delta| \leq u$, denotes the floating point arithmetic error. Therefore,

$$\text{fl} \left( \alpha_j \right) = \alpha_j + \xi_j,$$ \hspace{1cm} (4.2)

where

$$\xi_j = \alpha_j\delta + \sqrt{t + k_j + \alpha_j} \left( 1 + \delta \right).$$

Since $\alpha_j \leq \sqrt{n-j}\|Z\|$, we have

$$\xi_j \leq \alpha_j u + \left( \frac{|t|}{\alpha_j} \right) (1 + u) \leq \alpha_j u + \left( \frac{(n - j)\gamma_{n-j}\|Z\|^2}{\alpha_j} \right) (1 + u) \leq \varphi_j\|Z\|,$$ \hspace{1cm} (4.3)

where

$$\varphi_j := u\sqrt{n-j} + \frac{(n - j)\gamma_{n-j}\|Z\|}{\sqrt{k_0}}\|Z\|(1 + u).$$ \hspace{1cm} (4.4)
4.2 Error analysis of $\beta_j$

Let us consider

$$f_l(\beta_j) = \frac{\sinh(\alpha_j + \xi_j)}{\alpha_j + \xi_j}(1 + \delta).$$

By using the identity $\sinh x - \sinh y = 2 \sinh \frac{x-y}{2} \cosh \frac{x+y}{2}$, we obtain

$$f_l(\beta_j) = \left(\beta_j + \frac{\xi_j \sinh \frac{\xi_j}{2}}{\xi_j/2} \cosh \left(\frac{\alpha_j + \xi_j}{2} - \beta_j \xi_j\right)\right)(1 + \delta) = \beta_j + \varepsilon_j^1,$$

where the quantity $\varepsilon_j^1$ is

$$\varepsilon_j^1 = \beta_j \delta + \left(\frac{\xi_j \sinh \frac{\xi_j}{2}}{\xi_j/2} \cosh \left(\frac{\alpha_j + \xi_j}{2} - \beta_j \xi_j\right)\right)(1 + \delta). \quad (4.5)$$

From (4.1) and (4.3), we have

$$|\xi_j/\alpha_j| \leq u + \left(\frac{n - j}{k_0}\right)\frac{\gamma_{n-j}\|Z\|^2}{(1 + u)}. \quad (4.6)$$

We now consider the condition

$$\psi_j := u + \left(\frac{n - j}{k_0}\right)\frac{\gamma_{n-j}\|Z\|^2}{(1 + u)}. \quad (4.7)$$

that is clearly a consequence of

$$\|Z\|^2/k_0 < \frac{1}{1-u} - \frac{1}{1+n(1+u)} \quad \text{for } 1 \leq j \leq n-1,$$

which itself is obtained from condition (1.1b) for $\|Z\| < 1$. Hence from (4.7) and by increasing monotonicity of the map $x \mapsto \frac{x}{1+x}$ on the interval $(0, 1)$, we can write

$$\left|\frac{\xi_j}{\alpha_j + \xi_j}\right| \leq \frac{|\xi_j/\alpha_j|}{1 - |\xi_j/\alpha_j|} \leq \frac{\psi_j}{1 - \psi_j}. \quad (4.8)$$

The function

$$f(x) = \begin{cases} \sinh x, & x \neq 0, \\
1, & x = 0 \end{cases}$$

has the elementary property that $1 \leq f(a) \leq f(b)$ whenever $|a| \leq b$, from which and (4.5) we deduce $|\varepsilon_j^1| \leq \eta_j$, where

$$\eta_j := \frac{\sinh(\sqrt{n-j}\|Z\|)}{\sqrt{n-j}\|Z\|} \left(u + \frac{\psi_j}{1 - \psi_j}(1 + u)\right) + \frac{\psi_j \sinh(\psi_j\|Z\|/2)}{1 - \psi_j} \cosh \left(\sqrt{n-j}\|Z\| + \frac{\varphi_j\|Z\|^2}{2}\right)(1 + u). \quad (4.8)$$

Now by using (4.4) and (4.7) in (4.8), we get

$$|\varepsilon_j^1| \leq \sinh \left(u\sqrt{n-j}\|Z\|/2 + (n-j)^{\gamma_{n-j}\|Z\|^2/2}\right) + \frac{(n-j)^{\gamma_{n-j}\|Z\|^2/2}(1 + u)}{1 - u - (n-j)^{\gamma_{n-j}\|Z\|^2/2}(1 + u)} \times \cosh \left(1 + u/2\right) \sqrt{n-j}\|Z\| + (n-j)^{\gamma_{n-j}\|Z\|^2/2}(1 + u) \times \frac{u + (n-j)^{\gamma_{n-j}\|Z\|^2/2}(1 + u) + O(u^2)}{1 - u - (n-j)^{\gamma_{n-j}\|Z\|^2/2}(1 + u)}.$$
This shows that \( \varepsilon_j^1 \to 0 \) as \( \|Z\| \to 0 \), provided that \( \|Z\|_{k_0} \) admits an upper bound as suggested by the condition (1.1).

### 4.3 Error analysis of \( \lambda_j \)

However, for computing \( \lambda_j \), from (4.2) and definition of floating point arithmetic error \( \delta \) we have

\[
\sinh \left( \left( \frac{\alpha_j + \xi_j}{2} \right) (1 + \delta) \right) = \sinh \frac{\alpha_j}{2} \left( \cosh \frac{\alpha_j \delta + \xi_j (1 + \delta)}{2} + \coth \frac{\alpha_j}{2} \sinh \frac{\alpha_j \delta + \xi_j (1 + \delta)}{2} \right) = \sinh \frac{\alpha_j}{2} (1 + \varepsilon_j^2),
\]

where

\[
\varepsilon_j^2 = 2 \sinh^2 \frac{\alpha_j \delta + \xi_j (1 + \delta)}{4} + \coth \frac{\alpha_j}{2} \sinh \frac{\alpha_j \delta + \xi_j (1 + \delta)}{2}.
\]

Hence,

\[
\text{fl} (\lambda_j) = 2 \sinh^2 \frac{\alpha_j}{2} (1 + \varepsilon_j^2)^2 (1 + \theta_3).
\]

So,

\[
\text{fl} (\lambda_j) = \lambda_j (1 + \varepsilon_j^3),
\]

where

\[
\varepsilon_j^3 = \varepsilon_j^2 (\varepsilon_j^2 + 2) (1 + \theta_3) + \theta_3. \tag{4.9}
\]

Also, \( |\varepsilon_j^2| \leq \rho_j \), where

\[
\rho_j := 2 \sinh^2 \frac{\sqrt{n - j} \|Z\| u + \varphi_j (1 + u) \|Z\|}{4} + \coth \frac{\sqrt{k_0}}{2} \frac{\sqrt{n - j} \|Z\| u + \varphi_j (1 + u) \|Z\|}{2},
\]

\[
|\rho_j| \leq 2 \sinh^2 \left( \frac{1}{2} \frac{u \sqrt{n - j} \|Z\| + (n - j) \gamma_{n-j} \|Z\|^2}{4 \sqrt{k_0}} (1 + 2u) \right) + \coth \left( \frac{\sqrt{k_0}}{2} \right) \left[ \sinh \left( u \sqrt{n - j} \|Z\| + (n - j) \gamma_{n-j} \|Z\|^2 \right) + \frac{1}{2} \frac{(n - j) \gamma_{n-j} \|Z\|^2}{2} (1 + 2u) \right] + O(u^2).
\]

However,

\[
\coth \frac{\sqrt{k_0}}{2} = \cosh \frac{\sqrt{k_0}}{2} \frac{\sqrt{n - j} \|Z\| u + \varphi_j (1 + u) \|Z\|}{\sinh \frac{\sqrt{k_0}}{2}} \leq \cosh \frac{\sqrt{k_0}}{2},
\]

and therefore

\[
|\rho_j| \leq 2 \sinh^2 \left( \frac{1}{2} \frac{u \sqrt{n - j} \|Z\| + (n - j) \gamma_{n-j} \|Z\|^2}{4 \sqrt{k_0}} (1 + 2u) \right) + \cosh \left( \frac{\sqrt{k_0}}{2} \right) \left[ \sinh \left( u \sqrt{n - j} \|Z\| + (n - j) \gamma_{n-j} \|Z\|^2 \right) + \frac{1}{2} \frac{(n - j) \gamma_{n-j} \|Z\|^2}{2} (1 + 2u) \right] \times \left( 2u \sqrt{n - j} \|Z\| \frac{1}{\sqrt{k_0}} + (n - j) \gamma_{n-j} \|Z\|^2 \right) (1 + 2u) + O(u^2).
\]

Hence in this case \( \varepsilon_j^2 \to 0 \) as \( \|Z\| \to 0 \), provided that \( \frac{\|Z\|}{\sqrt{k_0}} \) admits an upper bound like the one suggested by condition (1.1) (note that \( \frac{\|Z\|}{\sqrt{k_0}} = \sqrt{\frac{\|Z\|}{k_0} \cdot \sqrt{\|Z\|}} \)). Moreover, by (4.9),

\[
|\varepsilon_j^3| \leq \alpha_j, \tag{4.10}
\]

where

\[
\alpha_j := (\rho_j^2 + 2\rho_j) (1 + \gamma_3) + \gamma_3, \tag{4.11}
\]

which implies \( \alpha_j \) is of order \( O(u) \) under condition (1.1).
4.4 Error analysis of \( \varphi(D) \)

Now, in calculating \( \varphi(D) \), we will have

\[
\text{fl}(\varphi(D)) = \left( \frac{\text{fl}(\beta_j)}{\text{fl}(\alpha_j)} \right)^{1 + \delta} = \left( \frac{\lambda_j(1 + \delta^j)}{(\alpha_j + \xi_j)^2(1 + \delta)} \right) = \varphi(D) + A,
\]

where

\[
A = \left( \begin{array}{c}
\beta_j \delta^j \\
\lambda_j \delta^j
\end{array} \right) \frac{\lambda_j \delta^j}{(1 + \xi_j/\alpha_j)^2 - 1}.
\]

From (4.6) and (4.10), we have

\[
|\sigma_j| \leq \frac{\alpha_j(1 + u) + u + 2\psi + \psi^2}{(1 - \psi)^2},
\]

(4.12)

Hence by (4.7) and (4.11), the right-hand side of (4.12) is of order \( O(u) \) and

\[
|\varphi(D)| \leq B,
\]

(4.13)

where

\[
B := \left( \begin{array}{cc}
\sinh \sqrt{n - j}\|Z\| & 2\sinh^2 \left( \sqrt{n - j}\|Z\|/2 \right) \\
2\sinh^2 \left( \sqrt{n - j}\|Z\|/2 \right) & (n - j)^2\|Z\|^2
\end{array} \right),
\]

(4.14)

\[
|A| \leq C,
\]

(4.15)

\[
C := \left( \begin{array}{cc}
\sinh \sqrt{n - j}\|Z\|/\eta_j & 2\sinh^2 \left( \sqrt{n - j}\|Z\|/2 \right) |\sigma_j| \\
2\sinh^2 \left( \sqrt{n - j}\|Z\|/2 \right) \eta_j & \sinh \sqrt{n - j}\|Z\|/\eta_j
\end{array} \right),
\]

(4.16)

4.5 Error analysis of \( w \)

At first, since \( v_k := \exp(z_{k,k})v_k \) in Algorithm 2, we have \( \bar{v}_{\text{old},i} = \phi_{\text{old},i}(1 + \theta^{(i)}_2) \), \( 1 \leq i \leq n \). Now, for computing \( w \), let \( v' = b_j^T v_{\text{old},i} \), hence,

\[
\text{fl}(v') = \sum_{i=j+1}^{n} \tau_{j,i} v_{\text{old},i} \left( 1 + \theta^{(i)}_2 \right) \left( 1 + \theta^{(i)}_{n-j} \right) = v' + \varepsilon_j^4,
\]

where

\[
\varepsilon_j^4 := \sum_{i=j+1}^{n} \tau_{j,i} v_{\text{old},i} \left( \theta^{(i)}_2 + \theta^{(i)}_{n-j} \theta^{(i)}_{n-j} \right).
\]

So from property (2.1), we have

\[
|\varepsilon_j^4| \leq (n - j)^2\|Z\|\|v\|\|e^z\|Z\|, \quad (4.17)
\]

Now from Algorithm 2 and (2.4),

\[
\text{fl}(w) = \left( \varphi(D) + A \right) \left( \varepsilon_j^4 \right) + u_j = w + q_j,
\]
Also, and hence $|u_j| \leq \gamma_2 |\varphi(D)| + A \left( |v_{\text{odd}}(1 + \theta_2)_{\varepsilon_j}^T \right) + u_j$, \hspace{1cm} (4.19)

Therefore, from (4.13), (4.15), (4.17), (4.18), and (4.19), we have $|q_j| \leq r_j$, where

$$r_j := \|v\| + \|Z\| \left( \gamma_2 \left( (\alpha_j + 1) + (w_2 + q_{j,2}) \right) + \frac{1 + \gamma_2}{\|Z\|} \right)$$

\hspace{1cm} (4.20)

Let us consider $\text{fl}(v_{\text{new}}) = \left( w_1 + q_{j,1} \right) a_j + \left( u_{1} + u_{2} \right) + (w_2 + q_{j,2}) e_1 + \left( u_{1} + u_{2} \right) = \text{v}_{\text{new}} + g_j,$

where

$$g_j := \|v\| + \|Z\| \left( \frac{\sinh (\sqrt{n - j} \|Z\|)}{\sqrt{n - j} \|Z\|} + \frac{2 \sinh^2 (\sqrt{n - j} \|Z\|) / 2}{\|Z\|} \right).$$

\hspace{1cm} (4.22)

Also, $w_2 = \lambda_j v_{\text{odd},1} + \beta_j b_j^T v_{\text{odd}}$, such that $|w_2| \leq l_2$, where

$$l_2 := \|v\| \left( \frac{\sinh (\sqrt{n - j} \|Z\|)}{\sqrt{n - j} \|Z\|} \right) (n - j) + 2 \sinh^2 (\|Z\| / 2).$$

\hspace{1cm} (4.23)

So from (4.20), (4.22), and (4.21), we have

$$|g_j| \leq \left( l_1 \|Z\| \gamma_2 + r_{j,1} \|Z\| \right) (1 + 2\gamma_2) + l_2 \gamma_2 + r_{j,2} (1 + 2\gamma_2) \cdot 1.$$
5 Sensibility of Algorithm 2

As it was shown in last section and from (4.23), the computation error of $\exp(Z)v$ depends on $\|Z\|$ which is calculated by Algorithm 1. We therefore consider $Z = Z + E$ in which the norm of $E$ is obtained from (3.3) and (3.4). Therefore, the relations (4.3) and (4.6) reduce to $|\xi_j| \leq \varphi_j$ and $|\xi_j| \leq \psi_j$ where

$$
\varphi_j = u \sqrt{n-j} \|Z\| + \frac{1}{\sqrt{k_0}} ((n-j) \gamma_{n-j} (\|Z\| + \|E\|)^2 + (n-j)(2\|E\||\|Z\| + \|E\|^2)) (1+u),
$$

$$
\psi_j = u \frac{1}{k_0} ((n-j) \gamma_{n-j} (\|Z\| + \|E\|)^2 + (n-j)(2\|E\||\|Z\| + \|E\|^2)) (1+u).
$$

After substituting $\varphi_j$ and $\psi_j$ in the appropriate formulas given in the last section, the error upper bounds are computed in terms of $\|E\|$. Accordingly, the estimate (4.17) becomes

$$
|\epsilon_j^4| \leq (n-j) \gamma_{n-j+2} \|Z\| \|v\| e^{\|Z\|} + (n-j)(1+ \gamma_{n-j+2}) \|E\| \|v\| e^{\|Z\|},
$$

that in turn provides the error bound given by (4.20) as follows:

$$
r_j := B \left( \frac{\gamma_2 \|v\|}{(n-j) \gamma_{n-j+2} \|Z\| \|v\| e^{\|Z\|} + (n-j)(1+ \gamma_{n-j+2}) \|E\| \|v\| e^{\|Z\|}} \right)
$$

$$
+ C \left( \frac{n-j \|Z\| \|v\| e^{\|Z\|} + (n-j) \gamma_{n-j+2} \|Z\| \|v\| e^{\|Z\|}}{1+ \gamma_2 \|v\| e^{\|Z\|}} \right)
$$

$$
+ \gamma_2 (B + C) \left( \frac{(n-j) \|Z\| \|v\| e^{\|Z\|} + (n-j) \gamma_{n-j+2} \|Z\| \|v\| e^{\|Z\|}}{1+ \gamma_2 \|v\| e^{\|Z\|}} \right).
$$

Now if we consider initial errors influencing the vectors $v(j : n) = v_{\text{old}} + v_{\text{new}}$, for $j \geq 2$, these vectors can be represented by $\tilde{v}(j : n) = v(j : n) + \text{error}(j)$, where error(1) = 0. Hence, the estimate (4.17) reduces to

$$
|\epsilon_j^4| \leq (n-j) \gamma_{n-j+2} \|Z\| \|v\| e^{\|Z\|} + (n-j)(1+ \gamma_{n-j+2}) \|E\| \|v\| e^{\|Z\|} + \|E\| \|\text{error}\|, \tag{5.1}
$$

in which the quantities $\|E\| \|\text{error}\|$, which are very small, can be neglected. Obviously, the upper bound given by (5.1) affects the variations of $r_j$ and errorj. Let “error” be the final computing error committed by both Algorithms 1 and 2. In Figure 1, we compared the dependence of $\|E\|$ and $\|\text{error}\|$ using $10 \times 10$ random matrices $hZ$ with $Z$ normalized, so that $\|Z\| = 1$, $h = 1, 2, \frac{1}{3}, \frac{1}{4}, \ldots, \frac{1}{10}$, and $10 \times 10$ random unit vector v. Figure 1 shows the stability of algorithms with respect to round off errors. Moreover, the statistical correlation coefficient between backward and forward errors is 0.9531.

![Figure 1: Backward and forward errors versus h for Algorithms 1 and 2.](image-url)
6 Conclusion

As we expected, the upper bound of $|\text{error}|$ (4.23) strongly depends on the input data $Z$ and $v$. More precisely, if $k_0 \gg u$, then this error bound can be controlled by $\max\{uc\|Z\|, \|v\|\}$. Finally, we have shown the stability of the GPD method under condition (1.1).

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