C OMMUNICATIONS

1. Introduction

In managerial experience we can find a problem of service centre optimal location. Location of those objects like manufactures, distributive and shopping centres, supply depots markedly affects the costs of material flows in creative logistic networks. The location of centers is so much complicated because there is not only one logistic chain but a whole distributive network.

Determination about a location or non-location of a service centre in some areas will affect the systems effectiveness for next several years. For finding the optimal solution it is possible to apply an exact method, but only for the known costs. When we solve location problems, for the most of them we have no real future costs, only their gross estimates. So it is necessary to deal with the approach of a location problem solving under uncertain costs.

This paper deals with a possible method of finding the optimal location of service centres under uncertain costs represented by fuzzy numbers.

2. Location problem formulation

A service centre can be set up only in some places from the finite set of possible locations, which requires standby costs. In the system are also costs of satisfying customer demands from some of located centres, which depend on quantity of requirements. The goal is to minimize complete costs of the system. So we have a difficult combinatorial problem of determination of a located service centres number.

There is securing freight traffic from one or more primary centres to customers in the distribution system. This freight traffic could be linear (without transshipments) or combined with transshipments in some centres called terminals, which are often warehouses or buffer stocks. The structure of distribution system is figured out by a set of primary centres, customers, terminals and flows of goods among them.

The location problem is a problem of optimal location of service centres on the given part of the transportation network.

The incapacitated location problem is conceived as follows:

The transportation network is given with customers in the nodes \( j \in J \) and localities \( i \in I \), in which it is possible to locate serve centres. Let’s also assume that also one centre located in the node from the set \( I \) is able to serve all customers (see Fig. 1). The task is to minimize complete costs, which include standby costs \( f_i \) paid for each location of the service centre in \( i \) and variable costs \( c_{ij} \) of demand satisfaction \( b_j \) of a customer \( j \) from the terminal \( i \).

The variable costs for satisfying demand \( b_j \) of customer \( j \in J \) is \( c_j = (e_1d_{si} + e_0d_{ij} + g_j)b_j \) consist of charges \( e_1 \) for import from the primary centre \( S \) to the terminal \( i \), costs \( g_j \) for transshipment in the transshipment \( i \) and charges \( e_0 \) for freight traffic from \( i \) to the customer \( j \). The haul between the primary centre \( S \) and the terminal \( i \) is \( d_{si} \) and between the terminal \( i \) and the customer \( j \) is \( d_{ij} \). The condition is that all the customers have to be served, or more precisely have to be assigned to some of the located terminals.

Having introduced 0–1 variable \( y_i \in \{0,1\} \) for each \( i \in I \), which models the decision if the terminal is located at \( i \) or not, and variable \( z_{ij} \in \{0,1\} \) for each pair \( i, j \in I, j \in J \), which assigns

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the customer \( j \) to the terminal location \( i \), we can set the following model of the complete cost minimization.

\[
\text{minimize } f(y, z) = \sum_{i \in I} f_i y_i + \sum_{i \in I} \sum_{j \in J} c_{ij} z_{ij}
\]

where \( c_{ij} = (e_i d_{ij} + e_d q_i) b_i \)  

subject to 

\[
\sum_{j \in J} z_{ij} = 1 \quad \text{for } j \in J \tag{1}
\]

\[
z_{ij} \leq y_i \quad \text{for } i \in I, j \in J \tag{2}
\]

\[
z_{ij} \geq 0 \quad \text{for } i \in I, j \in J \tag{3}
\]

\[
y_i \in [0, 1] \quad \text{for } i \in I. \tag{4}
\]

In the model above, the objective function (1) represents the complete costs of distribution system. Constraints (2) ensure that each customer demand has to be satisfied from exactly one terminal location, constraints (3) force placement of a terminal at location \( i \) whenever a customer is assigned to the terminal location \( i \), constraints (4) ensure the location of terminal in every locality from which demands of some customers are satisfied.

\section{3. Analysis of the existing approaches}

In strategic decision problems it is difficult to estimate future values of standby or/and variable costs. In this case, the estimation of future costs is inaccurate. Considering confidential variables, which model determination about (un)location of terminals, the resulting solution can be economically inefficient in the view of the future costs. For example, the growth of \( f_i \) or \( e_i \) creates a change of system structure of locations number and change of customers assignment (see Fig. 1). As a consequence, the estimation of expected costs by one numeric value is risky. Uncertain costs can be in that case described by an expert (see Fig. 4). So we solve the original task, but with a changed objective function describing uncertain costs:  

\[ F(x) = F_1(x), F_2(x), F_3(x) = \left( \sum_{i=1}^{n} q_i^1 x_i, \sum_{i=1}^{n} q_i^2 x_i, \sum_{i=1}^{n} q_i^3 x_i \right) \tag{7} \]

\section{4. Concept of location problem solving}

Another approach uses the theory of fuzzy sets, where the uncertain value \( q \) is described by a possible interval and membership function \( \mu_q \) (see Fig. 3) - it is a power of applicability of given element to \( q \). This membership function has a triangular form.

According to fuzzy arithmetic rules, fuzzy numbers can be mutually added, subtracted and multiplied and divided by a real number without loss of the triangular form. When the coefficients \( q_i \) of an objective function \( F = q_1 x_1 + q_2 x_2 + \ldots + q_n x_n \) of a linear programming problem are triangular fuzzy numbers, then the value of the objective function for a given set of variable values \( x = \langle x_1, x_2, \ldots, x_n \rangle \) is also a triangular fuzzy number:

\[
F(x) = (F^1(x), F^2(x), F^3(x)) = \left( \sum_{i=1}^{n} q_i^1 x_i, \sum_{i=1}^{n} q_i^2 x_i, \sum_{i=1}^{n} q_i^3 x_i \right) \tag{7}
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ability, is the fuzzy algorithm [3]. This approach is based on introducing the fuzzy set $F_s$, which expresses an assertion that “value of $F$ is small” with the membership function shown in Fig. 5, where $F^{\text{max}}$ and $F^{\text{min}}$ denote respectively minimal values of $F(x)$ and $F^*(x)$ over a set of feasible solutions of the problem.

In this approach we searched a feasible solution $x^*$, for which the membership function of the fuzzy set “$F(x)$ and $F^*$” obtains the maximal value $h$ (see Fig. 5).

![Fig. 5: Membership function of fuzzy sets $F(x)$, $F$ and their intersection $F(x)$ and $F_s$](image)

The maximal value $h$ of the membership function of the fuzzy set $F(x)$ and $F_s$, and for the given $x$ has to satisfy the following equality in the cases when $F^*(x) \leq F^{\text{max}}$ holds.

$$ F^*(x) + (F^*(x) - F^*(x))h = F^{\text{max}} - (F^{\text{max}} - F^{\text{min}})h \quad (9) $$

In other cases $h$ can be set to zero. For the former case we get

$$ h(x) = \frac{F^{\text{max}} - F^*(x)}{F^*(x) - F^*(x) + F^{\text{max}} - F^{\text{min}}} \quad (10) $$

and we seek for $x^*$ maximizing $h(x)$, which is a non-linear programming problem. The following numerical process [3] obtains an approximate solution of the problem.

1. Set $h$ to an initial positive value near zero.
2. Minimize the following objective function

$$ F^*(x) + (F^*(x) - F^*(x))h $$

over the set of feasible $x$ and denote $x^*(h)$ the associated optimal solution

3. Compute $h(x^*(h))$ according to (10).
4. If $|h - h(x^*(h))| < \epsilon$ then stop else set $h = h(x^*(h))$ and go to step 3.

As it can be noticed in Fig. 5 or derived from the expressions (9), the direct fuzzy approaches make use only of the left hand side of the membership function. It means, that part of fuzzy number from $F^2$ to $F^3$ is not taken into account (see Fig. 5).

To overcome this weakness of the above-mentioned fuzzy approaches, there is another fuzzy approach [1], which makes use of the membership function on its whole range. This approach resembles way in which random coefficients are processed, when their distribution of probability is known. In this probabilistic-like approach the interval $[0, 1]$ of possible values of the membership function is divided by real numbers of an arbitrary chosen finite set $H \subset [0, 1]$. Then for each fuzzy coefficient $c$ from the location model the values $c^1, c^2, \ldots, c^r$ are determined, so that the constraint $\mu_c(c_j) \in H$ holds for $k = 1, 2, \ldots, r$. This is possible concerning fact that the level of satisfaction of a fuzzy number centre is 1.

![Fig. 6: Level of satisfaction assignment to values of membership function](image)

Then we minimize the weighted sum function over the feasible solutions $D$.

$$ \min \left( \sum_{k=1}^{m+1} \left( \sum_{i=1}^{m+1} h_k x_i + \sum_{j=1}^{m+1} \sum_{k=1}^{m+1} c_k z_j \right) \right) $$

subject to $(y, z) \in D$

(11)

The operating name of this method is weights 2.

In the case, we don’t have more accurate information about uncertain costs, it means $h_k = 1$ for $\forall k = 1, 2, \ldots, m+1$ the method is named minisum 2.

If we use results of classical sensitivity analyses in the weighted sum function and find for which of those results is its value minimal, it is the method minisum 1. When we have nonzero weights, the method is named weights 1.

To compare and verify both approaches, there was implemented branch and bound method and built a software tool for sensitivity analysis and fuzzy processing of the location problem [1]. Functionality of the program was tested on 90 examples making use of the whole road system of Slovakia with 2906 dwelling places and 71 possible terminal locations. This way, in accordance with the primary source selection at 10 big towns of Slovakia 10 basic problems with predefined parameters $f, e_1, e_2$ were obtained. By three types of modification done independently with each of the three parameters, there 90 benchmarks were obtained, which were used in the experiments.

An average locations number for the method weights 2 is $9 \pm 0.9$.

Average objective function value for method weights 2 is $32294723 \pm 3190392$ Sk.
5. Conclusion

The fuzzy algorithm computes unique solution, which is not dependent on expert’s ability (like classical fuzzy method) and is resistant to future changes.

We have compared these approaches:

- sensitivity analysis and its usage by methods minisum 1, minisum 2, weights 1, weights 2,
- classical fuzzy method,
- fuzzy algorithm.

If a triangular fuzzy number describes the uncertain costs of location problem, both methods weights 2 and fuzzy algorithm are correct ways of finding the design of distribution system. These methods give similar results (see tables 1 and 2). There is difference 12% in number of placed terminals (average is 9 terminals).

We suggest to perform both approaches and resulting design take into account only if results of these methods differ slightly. In opposite case, we suggest to perform an additional cost analysis and make the fuzzy cost more precise.

Comparing the methods minisum 1 and minisum 2 with weights 2 is only the reference example, because these approaches don’t take weights into account.

One of the program outputs is a graphical representation of the solution, so the user can find out the stability of the optimal solution (see Fig. 7).

It is possible to change the values of a chosen parameter and also the method – analysis of sensitivity, classical fuzzy approach or fuzzy algorithm. The results are: object time, value of objective function, optimal number of terminals and their names and also associated customers.
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