On Increasing and Invariant Parking Sequences

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October 23, 2020

Abstract

The notion of parking sequences is a new generalization of parking functions introduced by Ehrenborg and Happ. In the parking process defining the classical parking functions, instead of each car only taking one parking space, we allow the cars to have different sizes and each takes up a number of adjacent parking spaces after a trailer $T$ parked on the first $z-1$ spots. A preference sequence in which all the cars are able to park is called a parking sequence. In this paper, we study increasing parking sequences and count them via bijections to lattice paths with right boundaries. Then we study two notions of invariance in parking sequences and present various characterizations and enumerative results.

1 Introduction

Classical parking functions were first introduced by Konheim and Weiss [4]. The original concept involves a linear parking lot with $n$ available spaces and $n$ labeled cars each with a pre-fixed parking preference. Cars enter one-by-one in order. Each car attempts to park in its preferred spot first. If a car finds its preferred spot occupied, it moves towards the exit and takes the next available spot. If there is no space available, the car exits without parking. A parking function of length $n$ is a preference sequence for the cars in which all cars are able to park (not necessarily in their preferred spaces). A formal definition for parking functions can be stated as follows.

Definition 1.1. Let $n$ be a positive integer and $\vec{a} = (a_1, a_2, ..., a_n)$ be a sequence of positive integers. Assume $a_{(1)} \leq a_{(2)} \leq \cdots \leq a_{(n)}$ is the non-decreasing rearrangement of $\vec{a}$. Then the sequence $\vec{a}$ is a parking function of length $n$ if and only if $a_{(i)} \leq i$ for all indices $i$. Equivalently, $\vec{a}$ is a parking function (of length $n$) if and only if for all $1 \leq i \leq n$,

$$\#\{j : a_j \leq i\} \geq i.$$  

\begin{equation}
\tag{1}
\end{equation}

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For example, the preference sequences $(1, 2, 3, 4)$, $(2, 1, 3, 4)$ or $(1, 2, 4, 1)$ are all parking functions, while $(2, 2, 4, 2)$ is not since it will have one car leave un-parked. It is well-known that the number of classical parking functions is $(n + 1)^{n-1}$. An elegant proof by Pollak (see [7]) uses a circle with $(n + 1)$ spots where the parking functions are the preference sequences that can park all $n$ cars without using the $(n + 1)$-th spot.

One well-known generalization of Definition 1.1 is the notion of vector parking functions, or $\vec{u}$-parking functions. Let $\vec{u}$ be a non-decreasing sequence $(u_1, u_2, u_3, \ldots)$ of positive integers. A $\vec{u}$-parking function of length $n$ is a sequence $(x_1, x_2, \ldots, x_n)$ of positive integers whose non-decreasing rearrangement $x(1) \leq x(2) \leq \cdots \leq x(n)$ satisfies $x(i) \leq u_i$. Equivalently, $(x_1, \ldots, x_n)$ is a $\vec{u}$-parking function if and only if for all $1 \leq i \leq n$,

$$\# \{j : x_j \leq u_i\} \geq i.$$  \hspace{1cm} (2)

Denote by $PF_n(\vec{u})$ the set of all $\vec{u}$-parking functions of length $n$. When $u_i = i$ we obtain the classical parking functions. When $u_i = a + b(i - 1)$ for some $a, b \in \mathbb{Z}^+$, it is known that the number of $\vec{u}$-parking functions is $a(a + bn)^{n-1}$; see e.g. [5].

The set of parking functions is a basic object lying in the center of combinatorics, with many connections and applications to other branches of mathematics and disciplines, such as storage problems in computer science, graph searching algorithms, interpolation theory, diagonal harmonics, and sandpile models. Because of their rich theories and applications, parking functions and their variations have been studied extensively in the literature. See [10] for a comprehensive survey on the combinatorial theory of parking functions.

There is a particular generalization of parking functions that was recently introduced by Ehrenborg and Happ [2, 3], called parking sequences. Again, there are $n$ cars trying to park in a linear parking lot. In this new model the car $C_i$ has length $y_i \in \mathbb{Z}_+$ for each $i = 1, 2, \ldots, n$. Call $\vec{y} = (y_1, y_2, \ldots, y_n)$ the length vector. There is a trailer $T$ of length $z - 1$ parked at the beginning of the street after which the $n$ cars park with car $C_i$ taking up $y_i$ adjacent parking spaces. Given a sequence $\mathbf{c} = (c_1, \ldots, c_n) \in \mathbb{Z}_+^n$ for $i = 1, 2, \ldots, n$ the cars enter the street in order, and car $C_i$ looks for the first empty spot $j \geq c_i$. If the spaces $j$ through $j + y_i - 1$ are all empty, then car $C_i$ parks in these spots. If $j$ does not exist or any of the spots $j + 1$ through $j + y_i - 1$ is already occupied, then there will be a collision and the car cannot park and has to leave the street. In this case, we say the parking fails.

**Definition 1.2.** Assume there are $z - 1 + \sum_{i=1}^n y_i$ parking spots along a street, with the first $z - 1$ occupied by a trailer. The sequence $\mathbf{c} = (c_1, \ldots, c_n)$ is called a parking sequence for $(\vec{y}, z)$ where $\vec{y} = (y_1, \ldots, y_n)$ if all $n$ cars can park without any collisions. We denote the set of all such parking sequences by $PS(\vec{y}; z)$.

For example, $\mathbf{c} = (3, 7, 5, 3)$ is a parking sequence for $(\vec{y}; z)$ where $\vec{y} = (1, 2, 2, 3)$ and $z = 4$. Figure 1 shows how the cars $C_1, \ldots, C_4$ would park along the street with the preference sequence $\mathbf{c}$. As given in [3], the number of parking sequences in $PS(\vec{y}; z)$ is

$$z \cdot (z + y_1 + n - 1) \cdot (z + y_1 + y_2 + n - 2) \cdots (z + y_1 + \cdots + y_{n-1} + 1).$$  \hspace{1cm} (3)

From (1) and (2) it is easy to see that any permutation of a $\vec{u}$-parking function is also a $\vec{u}$-parking function. This is however not true for parking sequences. Consider
as an example a one-way street with 4 spots and 2 cars with fixed length vector 
\( \vec{y} = (2, 2) \) and \( z = 1 \), (no trailer). Then whereas \( c = (1, 2) \) is a parking sequence for 
\( (\vec{y}; z) \), \( c' = (2, 1) \) is not. Thus, it is natural to ask which parking sequences \( c \) are 
invariant for \( (\vec{y}; z) \), that is, \( c \) is still a parking sequence for \( (\vec{y}; z) \) after the entries 
of \( c \) are permuted. Another question is which sequence remains a parking sequence 
when the cars enter the street in different orders. In other words, we want to know 
which preference sequence allows all the cars to park when the length vector \( \vec{y} \) is 
permuted to \( (y_\sigma(1), y_\sigma(2), ..., y_\sigma(n)) \) for an arbitrary permutation \( \sigma \) of length \( n \).

There are several basic notions and variations associated with parking functions 
and their generalizations. Usually, these notions lead to a study of special classes of 
parking functions that have some interesting property. One of such special classes 
is the set of increasing parking functions, which have non-decreasing entries and 
are counted by the ubiquitous Catalan numbers. It is only natural to ask for a 
generalization of this class in the set of parking sequences.

We study these questions in the present work. The rest of the paper is organized 
as follows. In section 2, we discuss increasing parking sequences and their connec-
tion to lattice paths. In section 3, we fix the length vector \( \vec{y} \) and characterize all 
permutation-invariant parking sequences when \( \vec{y} \) has some special characteristics.
Then in section 4, we characterize all parking sequences that remain valid for all 
permutations of \( \vec{y} \). We finish the paper with some closing remarks in section 5.

2 Increasing Parking Sequences

In this section, we consider all non-decreasing parking sequences for any given pair 
\( (\vec{y}; z) \). By convention, we write \( [x] = \{1, 2, ..., x\} \) and the interval \( [x, y] = \{x, x + 
1, ..., y\} \), where \( x, y \in \mathbb{Z}_+ \) and \( x \leq y \). Given any sequence \( b = (b_1, ..., b_n) \in \mathbb{Z}_n^+ \), let 
\( b_{\text{nc}} = (b(1), ..., b(n)) \) be the non-decreasing rearrangement of the entries of \( b \) 
and the \( i^{th} \) entry \( b(i) \) of \( b_{\text{nc}} \) is called the \( i^{th} \) order statistic of \( b \). Next, we define the 
final parking configuration for any given parking sequence.

**Definition 2.1.** Let \( c \in \text{PS}(\vec{y}; z) \). The final parking configuration of \( c \) is the 
arrangement of cars \( C_1, C_2, ..., C_n \) following the trailer \( T \) encoding their relative 
order on the street after they are done parking using the preference sequence \( c \).

For example, in Figure 1, the final parking configuration of \( c = (3, 7, 5, 3) \) is 
\( T, C_1, C_3, C_2, C_4 \).

The following inequalities analogous to (1) give a necessary condition for being 
a parking sequence.

**Lemma 2.1.** Suppose \( c = (c_1, ..., c_n) \in \text{PS}(\vec{y}; z) \) where \( \vec{y} = (y_1, ..., y_n) \). Then 
\( \#\{j \in [n] : c_j \leq z\} \geq 1 \) and for each \( 1 \leq t \leq n - 1 \),

\[
\#\{j : c_j \leq z + \sum_{i=0}^{t-1} y(n-i)\} \geq t + 1.
\]
Proof. We have \(\#\{j : c_j \leq z\} \geq 1\) because otherwise, there is no car whose preference is less than or equal to \(z\), thus no car parks on spot \(z\) and we obtain a contradiction. Suppose for some \(t \in [1, n-1]\), \(\#\{j : c_j \leq z + \sum_{i=0}^{t-1} y(n-i)\} \leq t\). Then in the final parking configuration on spots \([1, z + \sum_{i=0}^{t-1} y(n-i)]\), there are at most 1 trailer and \(t\) cars occupying a total of at most \(z - 1 + y(n) + y(n-1) + \cdots + y(n-t+1)\) spots. Thus, not all spots are used in the final parking configuration and this contradicts the fact that \(c \in \PS(y)\).

\[\square\]

Corollary 2.2. Let \(c = (c_1, ..., c_n) \in \PS(y)\) where \(y = (y_1, ..., y_n)\). Then \(c_{(1)} \leq y\) and for \(j = 2, ..., n\),

\[c_{(j)} \leq y + \sum_{i=0}^{j-2} y(n-i). \tag{5}\]

We note that the conditions of Lemma 2.1 are not sufficient. Using the same example as before, even though \(c = (1, 2)\) and \(c' = (2, 1)\) both satisfy (4) for \(y = (2, 2)\), \(c' \not\in \PS(y)\). In addition, for a parking sequence \(c \in \PS(y)\), its rearrangement \(c_{\text{nc}}\) is not necessarily a parking sequence. Consider the following example for \(y = (1, 1, 4)\) and \(z = 1\). \(c = (5, 6, 1)\) is in \(\PS(y)\) but \(c_{\text{nc}} = (1, 5, 6)\) is not.

Definition 2.2. A sequence \(c = (c_1, ..., c_n) \in \PS(y)\) is an increasing parking sequence for \((y; z)\) if \(c_1 \leq c_2 \leq \cdots \leq c_n\). We denote the set of all increasing parking sequences for \((y; z)\) by \(\text{IPS}(y; z)\).

When \(y = (1, 1, ..., 1)\) and \(z = 1\) (i.e. the trailer of length 0), Definition 2.2 leads to the classical increasing parking functions, which are counted by the Catalan numbers. It is well-known that classical increasing parking functions of length \(n\) are in one-to-one correspondence with Dyck paths of semilength \(n\), which are lattice paths from \((0, 0)\) to \((n, n)\) that lie strictly above and not touching the line \(y = x - 1\). This result can be generalized to increasing parking sequences.

First we show that an analog of (1) is enough to characterize increasing parking sequences.

Proposition 2.3. Let \((y; z) = (y_1, ..., y_n; z)\). Then \(c = (c_1, ..., c_n) \in \text{IPS}(y; z)\) if and only if \(c_1 \leq c_2 \leq \cdots \leq c_n\) and for all \(i \in [n]\),

\[c_i \leq z + \sum_{j=1}^{i-1} y_j. \tag{6}\]

Proof. Observe that if \(c\) is a non-decreasing preference sequence satisfying (6), then the cars will park in the final configuration \(T, C_1, ..., C_n\). Hence \(c\) is in \(\text{IPS}(y; z)\).

Conversely, for a non-decreasing sequence \(c\) that allows all the cars to park, we need to prove that it satisfies (6). First by Corollary 2.2, \(c_1 \leq z\). Thus, car \(C_1\) parks right after the trailer leaving no gap. By the rules of the parking process, if \(c_i \leq c_{i+1}\) and both cars \(C_i\) and \(C_{i+1}\) are able to park, then \(C_{i+1}\) will park after \(C_i\). Hence for a non-decreasing \(c \in \PS(y; z)\), the final parking configuration must be \(T, C_1, C_2, ..., C_n\). It follows that the first spot occupied by car \(C_i\) is \(z + y_1 + \cdots + y_{i-1}\), which must be larger than or equal to \(c_i\). \(\square\)
Proposition 2.3 allows us to enumerate increasing parking sequences for any given length vector \( \vec{y} \) and \( z \in \mathbb{Z}_+ \) using results in lattice path counting. Recall that a lattice path from \((0,0)\) to \((p,q)\) is a sequence of \( p \) east steps and \( q \) north steps. It can be represented by a sequence of non-decreasing integers \((x_1, x_2, \ldots, x_q)\) such that the north steps are at \((x_i, i - 1) \rightarrow (x_i, i)\), for \( i = 1, \ldots, q \). The lattice path is said to have strict right boundary \((b_1, b_2, \ldots, b_q)\) if \( 0 \leq x_i < b_i \) for all \( 1 \leq i \leq q \).

Let \( \text{LP}_{p,q}(b_1, b_2, \ldots, b_q) \) denote the set of all lattice paths from \((0,0)\) to \((p,q)\) with strict right boundary \((b_1, b_2, \ldots, b_q)\). Figure 2 shows an example of a lattice path \((2,3,3,7)\) from \((0,0)\) to \((8,4)\) with strict right boundary \(b = (3,4,5,8)\).

![Figure 2: A lattice path (2,3,3,7) with strict right boundary at (3,4,5,8).](image)

We can represent increasing parking sequences in terms of lattice paths with strict right boundary as follows: Let \((\vec{y}; z) = (y_1, \ldots, y_n; z)\) and \(M = z - 1 + y_1 + y_2 + \cdots + y_{n-1} + y_n\). By Proposition 2.3 there is a bijection from \(\text{IPS}(\vec{y}; z)\) to the set of lattice paths from \((0,0)\) to \((M,n)\) with strict right boundary \((z, z + y_1, z + y_1 + y_2, \ldots, z + y_1 + y_2 + \cdots + y_{n-1})\). The boundary is strict because in the lattice path, \(x_i\) can be 0 while in \(c \in \text{IPS}(\vec{y}; z)\), \(c_i \geq 1\). There are well-known determinant formulas to count the number of lattice paths with general boundaries, see, for example, Theorem 1 of [6, Chap.2], which leads to the following determinant formula.

**Corollary 2.4.** Suppose \( M = z - 1 + y_1 + y_2 + \cdots + y_{n-1} + y_n \). Then

\[
\#\text{IPS}(\vec{y}; z) = \#\text{LP}_{M,n}(z, z + y_1, z + y_1 + y_2, \ldots, z + y_1 + y_2 + \cdots + y_{n-1}) = \det \left[ b_i \right]_{1 \leq i,j \leq n}
\]

where \( b_1 = z \) and \( b_i = z + y_1 + y_2 + \cdots + y_{i-1} \) for \( i = 2, \ldots, n \).

For the special case that the length vector has constant entries, there are nicer closed formulæ for the determinant. Specifically, when \( \vec{y} = (k^n) = (k, k, \ldots, k) \) and \( M = z + kn - 1 \), \( \text{LP}_{M,n}(z, z + k, z + 2k, \ldots, z + (n - 1)k) \) is the set of lattice paths from \((0,0)\) to \((z + kn - 1, n)\) which never touch the line \( x = z + ky \). Using the formula (1.11) of [6, Chap.1], we have

**Corollary 2.5.** Suppose \((\vec{y}; z) = ((k^n); z)\) and \( M = z + kn - 1 \). Then

\[
\#\text{IPS}(\vec{y}; z) = \#\text{LP}_{M,n}(z, z + k, z + 2k, \ldots, z + (n - 1)k) = \frac{z}{z + n(k + 1)} \binom{z + n(k + 1)}{n}.
\]
This specializes to the Fuss-Catalan numbers when \( z = 1 \).

**Corollary 2.6.** Suppose \( \vec{y} = (k, k, \ldots, k) \in \mathbb{Z}_+^n \). Then

\[
\#\text{inv}(\vec{y}; 1) = \frac{1}{kn + 1} \binom{(k + 1)n}{n}.
\]

### 3 Invariance for Fixed Length Vectors

In this section, we study the first of two types of invariance for parking sequences. Fixing the length vector \( \vec{y} \in \mathbb{Z}_+^n \) and a positive integer \( z \), we investigate which parking sequence remains in the set \( \text{PS}(\vec{y}; z) \) after its entries are arbitrarily rearranged.

**Definition 3.1.** Fix \( \vec{y} = (y_1, \ldots, y_n) \) and \( z \in \mathbb{Z}_+ \). Let \( \vec{c} \in \text{PS}(\vec{y}; z) \). We say that \( \vec{c} \) is a permutation-invariant parking sequence for \( (\vec{y}; z) \) if for any rearrangement \( \vec{c}' \) of \( \vec{c} \), we have \( \vec{c}' \in \text{PS}(\vec{y}; z) \). We denote the set of all permutation-invariant parking sequences for \( (\vec{y}; z) \) by \( \text{PS}_{\text{inv}}(\vec{y}; z) \).

For example, for \( \vec{y} = (1, 2) \) and \( z = 1 \), we have \( \text{PS}(\vec{y}; 1) = \{(1, 1), (1, 2), (3, 1)\} \) and \( \text{PS}_{\text{inv}}(\vec{y}; z) = \{(1, 1)\} \). First, we describe a subset of the invariant parking sequences.

**Proposition 3.1.** For any \( \vec{c} = (c_1, \ldots, c_n) \in [z]^n \), we have \( \vec{c} \in \text{PS}_{\text{inv}}(\vec{y}; z) \).

**Proof.** For any preference sequence \( \vec{c} \), if \( c_i \leq z \) for all \( i \), then we obtain the final parking configuration \( T, C_1, C_2, \ldots, C_n \), which means \( \vec{c} \in \text{PS}(\vec{y}; z) \). Since the condition \( c_i \leq z \) for all \( i \) does not depend on the order of \( c_i \), we have that \( \vec{c} \) is permutation-invariant.

In general, \( \text{PS}_{\text{inv}}(\vec{y}; z) \) is larger than the set \([z]^n\), and the situation can be more complicated. The following two examples show that \( \text{PS}_{\text{inv}}(\vec{y}; z) \) depends not only on the relative order of the \( y_i \)'s, but also on the difference of \( y_i \)'s.

**Example 1.** Let \( \vec{y} = (y_1, y_2) \) and \( z = 1 \). If \( y_1 < y_2 \), then \( \text{PS}_{\text{inv}}(\vec{y}; 1) = \{(1, 1)\} \).

On the other hand, if \( y_1 \geq y_2 \), we have \( \text{PS}_{\text{inv}}(\vec{y}; 1) = \{(1, 1), (1, y_2 + 1), (y_2 + 1, 1)\} \).

**Example 2.** Suppose \( \vec{y} = (4, 3, 2) \) and \( \vec{t} = (4, 3, 1) \). It is easy to check that

\[
\text{PS}_{\text{inv}}(\vec{y}; 1) = \{(1, 1, 1), (1, 1, 4), (1, 4, 1), (4, 1, 1)\}
\]

and

\[
\text{PS}_{\text{inv}}(\vec{t}; 1) = \{(1, 1, 1), (1, 1, 4), (1, 4, 1), (4, 1, 1), (1, 1, 5), (1, 5, 1), (5, 1, 1)\}.
\]

Note that the relative orders for the vectors \( \vec{y} \) and \( \vec{t} \) are the same, (both have the pattern 321), but the invariant sets are not similar.

In the following we characterize the invariant set for some families of \( \vec{y} \). First, we consider the case where the length vector is strictly increasing. Next, we look at the case where \( \vec{y} \) is a constant sequence. Lastly, given \( a, b \in \mathbb{Z}_+ \), we consider the case where the length vector is of the form \( \vec{y} = (a, \ldots, a, b, \ldots, b) \), where \( a < b \).
3.1 Strictly increasing length vector

When \( \vec{y} \) is a strictly increasing sequence, we show that Proposition 3.1 gives all the permutation-invariant parking sequences.

**Theorem 3.2.** Let \( (\vec{y}; z) = (y_1, y_2, ..., y_n; z) \) where \( y_1 < y_2 < \cdots < y_n \). Then

\[
\text{PS}_{\text{inv}}(\vec{y}; z) = [z]^n.
\]

**Proof.** By Proposition 3.1, \([z]^n \subseteq \text{PS}_{\text{inv}}(\vec{y}; z)\). Conversely, suppose \( c = (c_1, c_2, ..., c_n) \) is a parking sequence for \((\vec{y}; z)\) with some \( c_i \not\in [z] \). We claim that \( c \) is not permutation-invariant. To see this, let \( x = \min\{c_i \mid c_i > z\} \). Then we can consider \( c_{\text{inc}} = (c_1, c_2, ..., c_r, x, c_{(r+2)}, ..., c_n) \), where \( c_1 \leq c_2 \leq \cdots \leq c_r \leq z < x \leq c_{(r+2)} \leq \cdots \leq c_n \) and \( r \geq 1 \) by Corollary 2.2. Then by Proposition 2.3, \( x \) satisfies the inequality: \( z < x \leq z + \sum_{i=1}^r y_i \). Thus, we can choose the maximum \( s \) such that \( x > z + \sum_{i=1}^s y_i \), where \( 0 \leq s < r \). Consider the preference

\[
c' = (c_1, c_2, ..., c_s, x, c_{(s+1)}, ..., c_r, c_{(r+2)}, ..., c_n)
\]

We try to park according to \( c' \). Clearly, the first \( s \) cars park in order after the trailer \( T \) without any gaps in between them. Then the car \( C_{s+1} \) has preference \( x \) and parks after car \( C_s \) with \( h \) unoccupied spots in between \( C_s \) and \( C_{s+1} \), where \( h = x - (z + \sum_{i=1}^s y_i) \geq 1 \) and \( h \leq y_{s+1} \) by the maximality of \( s \). Among the un-parked cars \( C_{s+2}, ..., C_n \), the minimal length is \( y_{s+2} \), where \( y_{s+2} > y_{s+1} \geq h \). Thus no car can fill in these \( h \) unoccupied spots. It follows that \( c' \not\in \text{PS}(\vec{y}; z) \), and hence \( c \not\in \text{PS}_{\text{inv}}(\vec{y}; z) \).

**Corollary 3.3.** Let \( (\vec{y}; z) = (y_1, y_2, ..., y_n; z) \) where \( y_1 < y_2 < \cdots < y_n \). Then

\[
\#\text{PS}_{\text{inv}}(\vec{y}; z) = z^n.
\]

3.2 Constant length vector

In this subsection, we investigate the case where \( \vec{y} \) is of the form \((k^n) = (k, k, ..., k)\).

**Theorem 3.4.** Suppose \( (\vec{y}; z) = ((k^n); z) \) where \( k \in \mathbb{Z}_+ \) and \( k > 1 \). Then \( \text{PS}_{\text{inv}}(\vec{y}; z) \) is the set of all sequences \((c_1, ..., c_n)\) such that for each \( 1 \leq i \leq n \),

(i) \( c_{(i)} \leq z + (i-1)k \), and

(ii) \( c_i \in \{1, 2, ..., z, z+k, z+2k, ..., z+(n-1)k\} \).

**Proof.** Let \( c = (c_1, ..., c_n) \) be a sequence that for each \( 1 \leq i \leq n \), \( c_{(i)} \leq z + (i-1)k \), and \( c_i \leq z \) or \( c_i = z + sk \) for some \( s = 0, 1, ..., n-1 \). We claim that \( c \in \text{PS}(\vec{y}; z) \).

Since these conditions are independent of the arrangement of the terms \( c_i ' s \), this implies \( c \) is permutation-invariant.

We attempt to park using \( c \). First, \( C_1 \) either parks right after the trailer if \( c_1 \leq z \), or on spots \( (c_1, c_1 + k - 1) \) if \( c_1 = z + sk \) for some \( s \geq 0 \). We assume for our inductive hypothesis, that the first \( r \) cars are parked already, (where \( 1 \leq r \leq n-1 \)).

At this stage in the parking process, any car already parked on the street occupies spots of the form \([z + ks, z + k(s+1) - 1]\) where \( s \in \{0, 1, ..., n-1\} \). This means that for any maximal interval of unoccupied spots, the length is a multiple of \( k \) and the interval starts at \( z + km \) for some \( m \in \{0, 1, ..., n-1\} \).

Now car \( C_{r+1} \) comes in with preference \( c_{r+1} \leq z \) or \( c_{r+1} = z + kl \). If \( c_{r+1} \leq z \), clearly \( C_{r+1} \) can park. If \( c_{r+1} = z + kl \) for some \( l \geq 1 \), then there are two possibilities:
• if spot \((z + kl)\) is empty, then \(C_{r+1}\) parks on spots \([z + kl, z + k(l + 1) - 1]\).

• if spot \((z + kl)\) is non-empty, then \(C_{r+1}\) drives forward to park in the first open interval ahead. Such an open interval must exist. Otherwise, assume that the last open spot after \(C_r\) parked is \(x\). By the inductive hypothesis, \(x = z + sk - 1\) for some \(s < l\). So all the spots from \(x + 1\) to the end of the street \((z + nk - 1)\) are occupied, by \(n - s\) cars. These \(n - s\) cars, as well as \(C_{r+1}\), all have preference greater than \(x\). In other words, there are at least \(n - s + 1\) cars having preference \(c_i \geq z + sk\), which implies \(c_{(s)} \geq z + sk\), a contradiction.

This exhausts all possible cases for \(C_{r+1}\). Thus, by induction, all cars can park and \(c \in \text{PS}(\bar{y}; z)\).

Conversely, suppose for a contradiction that there is a parking sequence \(c \in \text{PS}_{\text{inv}}(\bar{y}; z)\) not satisfying Condition (ii). (By Corollary 2.2 Condition (i) holds for all \(c \in \text{PS}(\bar{y}; z)\).) Then there is some \(j \in [n]\) such that \(c_j = z + sk + t\) for some \(s \in \{0, 1, ..., n - 1\}\) and \(1 \leq t < k\). Consider the following rearrangement of \(c\) given by \(c' = (c_j, c_1, c_2, ..., c_{j-1}, c_{j+1}, ..., c_n)\). By our assumption, \(c' \in \text{PS}(\bar{y}; z)\). We attempt to park using this preference. First, \(C_1\) parks on \([z + sk + t, z + (s + 1)k + t - 1]\). However, between the trailer and \(C_1\), there is now an unoccupied interval of \((z + sk + t) - z = sk + t\) spots, which is clearly nonempty and not a multiple of \(k\). Thus, no matter what preferences the remaining cars have, it is impossible to park all cars on this street. This yields a contradiction to our assumption.

Recall that a \(\bar{u}\)-parking function of length \(n\) is an integer sequence \((x_1, x_2, ..., x_n)\) satisfying \(1 \leq x_{(i)} \leq u_i\). We can use the results of vector parking functions to enumerate the number of sequences described in Theorem 3.4.

**Corollary 3.5.** Let \((\bar{y}; z) = ((k^n); z)\) with \(k \geq 2\). Then

\[
\# \text{PS}_{\text{inv}}(\bar{y}; z) = z(n + z)^{n-1}.
\]

**Proof.** For any \(c = (c_1, ..., c_n) \in \text{PS}_{\text{inv}}(\bar{y}; z)\) let \(f(c) = c'\), where \(c'\) is the sequence whose entries are given by

\[
c'_i = \begin{cases} 
c_i, & \text{if } 1 \leq c_i \leq z \\
z + s, & \text{if } c_i = z + sk.
\end{cases}
\]

The condition \(c_{(i)} \leq z + (i - 1)k\) implies \(c'_{(i)} \leq z + i - 1\), hence \(c'\) is a vector parking function associated to the vector \(\bar{u} = (z, z + 1, ..., z + n - 1)\), and \(f\) is a map from \(\text{PS}_{\text{inv}}(\bar{y}; z)\) to \(\text{PF}_n(\bar{u})\). It is clear that \(f\) is a bijection since the map can be easily inverted. By [5, Corollary 5.5], the number of \(\bar{u}\)-parking functions is \(z(z + n)^{n-1}\). \(\square\)

**Remark.** Note that Theorem 3.4 and Corollary 3.5 hold trivially for \(k = 1\), where the condition (ii) in Theorem 3.4 is redundant. In fact, when \(k = 1\) or \(n = 1\), \(\text{PS}_{\text{inv}}((k^n); z)\) is the same as the set \(\text{PS}((k^n); z)\), which is exactly the set of \(\bar{u}\)-parking functions associated to \(\bar{u} = (z, z + 1, ..., z + n - 1)\). For \(k, n \geq 2\), \(\text{PS}_{\text{inv}}((k^n); z)\) is a proper subset of \(\text{PS}((k^n); z)\).
3.3 Length vector $\vec{y} = (a, ..., a, b, ..., b)$ where $a < b$

Let $n \geq 2$ and $z, a, b, r$ be positive integers with $a < b$ and $1 \leq r < n$. In this subsection we fix $\vec{y} = (a, a, ..., a, b, ..., b) = (a^r, b^{n-r})$, i.e. the first $r$ cars are of size $a$ and the remaining $n - r$ cars are of size $b$. First we prove a couple of lemmas that characterize the set of permutation-invariant parking sequences for $(\vec{y}; z)$. In the following we will refer to any car of size $a$ (respectively, size $b$) as an $A$-car (respectively, $B$-car).

**Lemma 3.6.** Assume $c \in \text{PS}_{inv}(\vec{y}; z)$. Then in the final parking configuration of $c$, all $A$-cars park in $[z, z + ra - 1]$.

**Proof.** Suppose not. Then there is some $c \in \text{PS}_{inv}(\vec{y}; z)$ with which at least one $A$-car is not parked in the interval $[z, z + ra - 1]$ in the final parking configuration $\mathcal{F}$. Removing the trailer $T$ and all the $A$-cars from $\mathcal{F}$, we are left with some blocks of consecutive spots occupied by $B$-cars. Assume the first such block $L$ consists of $lb$ spots, where $1 \leq l \leq n - r$. Let $C_j$ be the last $A$-car in the configuration $\mathcal{F}$. Then $C_j$ occupies some spots in $[z + ra, z + ra + (n - r)b - 1]$, and no other $A$-car has checked the spots $C_j$ occupied in the parking process. In addition, let $C_k$ be the first $B$-car in $\mathcal{F}$. Then $j \leq r < k$ and $C_k$ parks before $C_j$ in $\mathcal{F}$. Let $c'$ be any permutation of $c$ in which the first $r$ terms are $c_1, ..., c_{j-1}, c_{j+1}, ..., c_r, c_k$, and the $(r + 1)$-th term is $c_j$. Note that these terms appear in the same order in both $c$ and $c'$.

**Case 1:** Assume in $\mathcal{F}$ there are some other $A$-cars parked between $C_k$ and $C_j$. First let the $A$-cars park according to the preference $c'$. It is easy to see that these $A$-cars occupy the same spots as in $\mathcal{F}$ except that

(a) there is no $A$-car taking the same spots as $C_j$ in $\mathcal{F}$, and

(b) In the $b$ spots originally occupied by $C_k$ in $\mathcal{F}$, there is now an $A$-car taking the first $a$ spots, leaving $(b - a)$ of these spots unoccupied.

Hence after all the $A$-cars are parked, the first block of consecutive open spots has size $lb - a$, which is not a multiple of $b$. Thus it is impossible for all the remaining $B$-cars to park and hence $c' \notin \text{PS}(\vec{y}; z)$.

**Case 2:** There is no $A$-car parked between $C_k$ and $C_j$ in $\mathcal{F}$. Then $\mathcal{F}$ is of the form $A \cdots AB \cdots BAB \cdots B$, where there are $r - 1$ $A$-cars before the first $B$-car $C_k$, and $c_j = z + (r - 1)a + lb$. Let the cars park according to $c'$. Then the $A$-cars all park successfully, occupying the spots $[z, z + ra - 1]$, and the first $B$-car occupies spots $[c_j, c_j + b - 1]$. Now there are $c_j - 1 - (z - 1 + ra) = lb - a$ spots between the last $A$-car and the first $B$-car; these spots cannot be filled by other $B$-cars. Hence $c' \notin \text{PS}(\vec{y}; z)$.

In both cases we have a permutation of $c$ that is not in $\text{PS}(\vec{y}; z)$, contradicting the assumption that $c \in \text{PS}_{inv}(\vec{y}; z)$.

**Lemma 3.7.** If $(c_1, c_2, ..., c_n) \in \text{PS}_{inv}(\vec{y}; z)$, then $c_i \leq z + (r - 1)a$ for all $1 \leq i \leq n$.

**Proof.** Suppose not. Take any permutation of $c$ starting with $\max\{c_i : i \in [n]\}$ and we contradict the conclusion of Lemma 3.6.

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Proposition 3.8. For \( c \in \text{PS}_{\text{inv}}(\vec{y}; z) \), let \( c(1) \leq c(2) \leq \cdots \leq c(n) \) be the order statistics of \( c \). Then \( \text{PS}_{\text{inv}}(\vec{y}; z) \) is exactly the set of all sequences \( c \) whose order statistics satisfy \( c(i) \leq z \) for each \( 1 \leq i \leq n - r + 1 \) and \( c(n-r+j) \in \{1, \ldots, z+a, z+2a, \ldots, z+(j-1)a\} \) for each \( 2 \leq j \leq r \).

Proof. By Lemmas 3.6 and 3.7, if \( c \in \text{PS}_{\text{inv}}(\vec{y}; z) \), then any \( r \)-term subsequence of \( c \), say \( (c_i, c_{i+1}, \ldots, c_{i+r}) \), parks all \( r \)-cars in \([z, z+ra-1] \) and hence \( c_i \leq z+(r-1)a \) for all \( i = 1, \ldots, n \). Furthermore, if we consider the last \( r \) terms of the order statistics of \( c \), this means \( (c(n-r+1), c(n-r+2), \ldots, c(n)) \in \text{PS}_{\text{inv}}((a, a, \ldots, a); z) \). By Theorem 3.4, we obtain \( c(n-r+j) \in \{1, \ldots, z+a, z+2a, \ldots, z+(j-1)a\} \) for each \( 1 \leq j \leq r \). Finally by the order statistics, \( c(i) \leq c(n-r+1) \leq z \) for each \( 1 \leq i \leq n-r \).

Conversely, we claim that any \( c \) satisfying the inequalities in Lemma 3.8 is in \( \text{PS}_{\text{inv}}(\vec{y}; z) \). To see this, consider first the \( A \)-cars with preferences \((c_1, \ldots, c_r)\). We have \( c_i \in \{1, 2, \ldots, z+a, z+2a, \ldots, z+(r-1)a\} \) for all \( 1 \leq i \leq r \), and the order statistics of these \( r \) terms are no more than \((z, z+a, \ldots, z+(r-1)a)\) (coordinate-wise). By Theorem 3.4, \((c_1, \ldots, c_r)\) is a parking sequence for \((a^r); z)\). Hence all \( A \)-cars must park on \([z, z+ra-1]\). Next, consider the \( B \)-cars. Since \( c_i \leq z+(r-1)a \) and all \( A \)-cars are parked without any unoccupied spots on \([z, z+ra-1]\), then all \( B \)-cars park in increasing order after the \( A \)-cars. In other words, the final parking configuration is \( T, C'_1, \ldots, C'_r, C_{r+1}, \ldots, C_n \) where \( C'_1, \ldots, C'_r \) is some rearrangement of the \( A \)-cars. 

Combining lemmas 3.6, 3.7 and Proposition 3.8, we prove the following result.

Theorem 3.9. Let \( n \geq 2 \) and \( z, a, b, r \) be positive integers with \( a < b \) and \( 1 \leq r < n \). Assume \( \vec{y} = (a^r, b^{n-r}) \). Let \( \text{PF}_n(\vec{u}) \) be the set of \( \vec{u} \)-parking functions of length \( n \) where \( \vec{u} = (z, z, \ldots, z, z+1, z+2, \ldots, z+r-1) \). Then there is a bijection between

\[ \text{PS}_{\text{inv}}(\vec{y}; z) \text{ and } \text{PF}_n(\vec{u}). \]

Proof. Let \( \vec{u} = (u_1, u_2, \ldots, u_n) = (z, z, \ldots, z, z+1, z+2, \ldots, z+r-1) \), where there are \( n-r+1 \) copies of \( z \) in the vector \( \vec{u} \). Consider the map \( \gamma_a : \text{PS}_{\text{inv}}(\vec{y}; z) \rightarrow \text{PF}_n(\vec{u}) \) defined as follows.

\[ \gamma_a : (c_1, \ldots, c_n) \mapsto (c'_1, \ldots, c'_n) = c' \]

where for all \( 1 \leq j \leq n \)

\[ c'_j = \begin{cases} c_j, & \text{if } c_j \leq z \\ z+s, & \text{if } c_j = z+sa. \end{cases} \]

The map \( \gamma_a \) is well-defined since the sequence \( c' \) has order statistics satisfying \( 1 \leq c'(i) \leq u_i \) for each \( i = 1, 2, \ldots, n \). Thus \( c' \in \text{PF}_n(\vec{u}) \). Clearly the map \( \gamma_a \) is invertible, hence \( \gamma_a \) is a bijection.

Corollary 3.10. Let \( \vec{y} \) and \( \vec{u} \) be as in Theorem 3.9. Then

\[ \#\text{PS}_{\text{inv}}(\vec{y}; z) = \#\text{PF}_n(\vec{u}) = \sum_{j=0}^{r-1} \binom{n}{j} (r-j)^{r-1} z^{n-j}. \]

In particular, when \( r = 1 \), \( \#\text{PS}_{\text{inv}}(a, b, b, \ldots, b; z) = z^n. \)
Proof. The result follows from Theorem 3.9 and [[9], Theorem 3].

It is natural to ask what happens for the length vector \( \vec{y} = (a^r, b^{n-r}) \) with \( a > b \). Unlike in the preceding subsection, the number of sequences in \( \text{PS}_{\text{inv}}((a^r, b^{n-r}); z) \) with \( a > b \) depends not only on \( r \) and \( z \), but also on the values of \( a \) and \( b \). We do not have a complete characterization yet, except some minor results for special cases.

4 Invariance for the Set of Car Lengths

In this section, we study another type of invariance. Given a fixed set of cars of various lengths and a one-way street whose length is equal to the sum of the car lengths and a trailer’s length \( z-1 \), we consider the parking sequences for which all \( n \) cars can park on the street irrespective of the order in which they enter the street. Denote by \( \mathcal{S}_n \) the set of all permutations on \( n \) letters. For a vector \( \vec{y} \) and \( \sigma \) in \( \mathcal{S}_n \), let \( \sigma(\vec{y}) = (y_{\sigma(1)}, \ldots, y_{\sigma(n)}) \).

Definition 4.1. Let \( c = (c_1, \ldots, c_n) \) and \( \vec{y} = (y_1, \ldots, y_n) \). Then \( c \) is a strong parking sequence for \( (\vec{y}; z) \) if and only if
\[
\forall \sigma \in \mathcal{S}_n \quad \sigma(\vec{y}) \in \text{PS}(\sigma(\vec{y}); z).
\]

We will denote the set of all strong parking sequences for \( (\vec{y}; z) \) by \( \text{SPS}\{\vec{y}; z\} \), or equivalently, \( \text{SPS}\{\vec{y}_{\text{inc}}; z\} \).

Example 3. For the case \( n = 2 \), let \( a, b \in \mathbb{Z}_+ \) with \( a < b \). It is easy to see that
\[
\text{PS}((a, b); z) = [z] \times [z + a] \cup \{(c_1, c_2) : c_1 = z + b, 1 \leq c_2 \leq z\},
\]
\[
\text{PS}((b, a); z) = [z] \times [z + b] \cup \{(c_1, c_2) : c_1 = z + a, 1 \leq c_2 \leq z\}.
\]

This gives
\[
\text{SPS}\{(a, b); z\} = \text{PS}((a, b); z) \cap \text{PS}((b, a); z) = [z] \times [z + a].
\]

Note that \( \text{SPS}\{(a, b); z\} \) is exactly the set of all preferences \( c \in \text{PS}((a, b); z) \) that yields the final parking configuration \( T, C_1, C_2 \).

By Ehrenborg and Happ’s result (3), we know that if \( \vec{y} = (k^n) \), then
\[
\#\text{SPS}\{\vec{y}; z\} = \#\text{PS}(\vec{y}; z) = z \cdot \prod_{i=1}^{n-1} (z + ik + n - i).
\]

In the following we consider the case where \( \vec{y} \) does not have constant entries.

Definition 4.2. We say that \( c \in \text{PS}(\vec{y}; z) \) parks \( \vec{y} \) in the standard order if the final parking configuration of \( c \) is given by \( T, C_1, C_2, \ldots, C_n \).

The following lemma is easily proved by induction.

Lemma 4.1. Let \( c = (c_1, c_2, \ldots, c_n) \in \text{PS}(\vec{y}; z) \). Then \( c \) parks \( \vec{y} \) in the standard order if and only if
\[
c_k \leq z + y_1 + \cdots + y_{k-1} \quad \text{for all} \quad k \in [n].
\]
The following result characterizes strong parking sequences for any set of \( n \geq 2 \) cars with a given length vector \( \{y_1, y_2, \ldots, y_n\} \) and a trailer \( T \) of length \( z - 1 \).

**Theorem 4.2.** Let \( n \geq 2 \). Assume that \( \vec{y} = (y_1, \ldots, y_n) \) is not a constant sequence. Then \( \mathbf{c} \) is a strong parking sequence for \( \{\vec{y}; z\} \) if and only if \( \mathbf{c} \) parks \( \vec{y}_{\text{inc}} = (y_1; y_2, \ldots, y(n)) \) in the standard order.

**Proof.** Suppose \( \mathbf{c} \) parks \( \vec{y}_{\text{inc}} \) in the standard order. We need to check that \( \mathbf{c} \) is a parking sequence for \( (\sigma(\vec{y}); z) \) for every \( \sigma \in \mathfrak{S}_n \). This follows from Lemma 4.1 and the fact that \( y_{(1)} + y_{(2)} + \cdots + y_{(i)} \leq y_{\sigma(1)} + y_{\sigma(2)} + \cdots + y_{\sigma(i)} \) for any \( \sigma \in \mathfrak{S}_n \) and \( i \in [n] \).

Conversely, let \( \mathbf{c} \) be a parking sequence for \( (\vec{y}_{\text{inc}}; z) \) that does not park \( \vec{y}_{\text{inc}} \) in the standard order. We will construct a permutation \( \sigma \) such that for a sequence of cars with length vector \( \sigma(\vec{y}_{\text{inc}}), \mathbf{c} \not\in \text{PS}(\sigma(\vec{y}_{\text{inc}}); z) \). In the following, let \( C_i \) represent a car of length \( y_{(i)} \), as listed in the table below. Let \( \mathcal{F} \) be the final parking configuration of \( \mathbf{c} \) when we park the cars \( C_1, \ldots, C_n \). In \( \mathbf{c} \), let \( k_1 \) be the minimal index \( k \) such that \( c_k > z + y_{(1)} + y_{(2)} + \cdots + y_{(k-1)} \). Then in \( \mathcal{F} \) the trailer is followed by \( C_1, \ldots, C_{k_1-1} \) with no gap, but there is a gap between \( C_{k_1-1} \) and \( C_{k_1} \). Let \( C_t \) be the last car that parks right before \( C_{k_1} \) in \( \mathcal{F} \). Clearly \( t > k_1 \).

| Car | \( C_1 \) | \( C_2 \) | \ldots | \( C_{k_1} \) | \( C_t \) | \ldots | \( C_{n-1} \) | \( C_n \) |
|-----|------|------|-----|------|------|-----|------|------|
| \( y_{(1)} \) | \( y_{(2)} \) | \ldots | \( y_{(k_1)} \) | \( y_{(t)} \) | \ldots | \( y_{(n-1)} \) | \( y_{(n)} \) |

**Case 1.** Assume \( y_{(k_1)} < y_{(t)} \). Let \( \sigma_1 \) be the transposition \( (k_1 \leftrightarrow t) \). For each \( i \in [n] \) let \( D_i \) represent a car of length \( y_{\sigma_1(i)} \), as shown below.

| \( \sigma_1 \) | Car | \( D_1 \) | \( D_2 \) | \ldots | \( D_{k_1} \) | \( D_t \) | \ldots | \( D_{n-1} \) | \( D_n \) |
|-----|------|------|-----|------|------|-----|-----|------|------|
| \( y_{(1)} \) | \( y_{(2)} \) | \ldots | \( y_{(k_1)} \) | \( y_{(t)} \) | \ldots | \( y_{(n-1)} \) | \( y_{(n)} \) |

We park cars \( D_1, \ldots, D_n \) using the preference sequence \( \mathbf{c} \). If \( \mathbf{c} \in \text{PS}(\sigma_1(\vec{y}_{\text{inc}}); z) \), then \( D_1, \ldots, D_t \) are able to park and

(a) \( D_1, D_2, \ldots, D_{k_1-1} \) have the same lengths and preferences as \( C_1, C_2, \ldots, C_{k_1-1} \). Hence they park in order right after the trailer with no gaps.

(b) \( D_{k_1} \) is longer than \( C_{k_1} \) and occupies spots in \([c_{k_1}, c_{k_1} + y_{(t)} - 1]\).

(c) Any car \( D_i \) for \( i \in \{k_1 + 1, \ldots, t - 1\} \) has the same preference as \( C_i \) so it parks either before \( D_{k_1} \) and in the same spots as \( C_i \) in \( \mathcal{F} \), or parks after \( D_{k_1} \).

(d) \( D_t \) takes the first \( y_{(k_1)} \) spots of the ones occupied by \( C_t \) in \( \mathcal{F} \).

After parking \( D_1, \ldots, D_t \), there are \( y_{(t)} - y_{(k_1)} \) unused spots between cars \( D_t \) and \( D_{k_1} \). Any car trying to park after \( D_t \) has length \( y_{(t)} > y_{(t)} - y_{(k_1)} \). So the spots between \( D_t \) and \( D_{k_1} \) cannot be filled and hence \( \mathbf{c} \not\in \text{PS}(\sigma_1(\vec{y}_{\text{inc}}), z) \).

**Case 2.** Assume \( y_{(k_1)} = y_{(t)} \). Then since \( \vec{y}_{\text{inc}} \) is not a constant sequence, either \( y_{(1)} < y_{(k_1)} \) or \( y_{(1)} < y_{(n)} \).

**Case 2a:** Assume \( y_{(1)} < y_{(n)} \). Let \( \sigma_2 \) be the transposition \( (t \leftrightarrow n) \) and \( E_i \) be a car of length \( \sigma_2(i) \) for each \( i \in [n] \).
We park cars \(E_1, \ldots, E_n\) using the preference sequence \(c\). The cars \(E_1, \ldots, E_{t-1}\) take the same spots as \(C_1, \ldots, C_{t-1}\) in \(F\). Next, car \(E_t\) tries to park in the spots \(C_t\) occupies, at the interval \([c_{k_t} - y(t), c_{k_t} - 1]\), where the spot \(c_{k_t}\) is already occupied by \(E_{k_t}\). But \(E_t\) has length \(y(n) > y(t)\) and hence cannot fit. Therefore, \(c \notin PS(\sigma_2(y_{inc}); z)\).

**Case 2b:** If \(y(k_1) = \ldots = y(t) = \ldots = y(n) = b\), then we must have \(k_1 > 1\) and \(y(1) < y(k_1)\). In the final configuration \(F\), at the time car \(C_{k_1}\) is parked, the lengths of all the intervals of consecutive empty spots left are multiples of \(b\). Let \(\sigma_3\) be the transposition \((1 \leftrightarrow k_1)\) and \(F_1\) be a car of length \(\sigma_3(i)\) for each \(i \in [n]\).

| \(\sigma_3\) | Car | Car Length | \(F_1\) | \(F_2\) | \(\ldots\) | \(F_{k_1}\) | \(\ldots\) | \(F_t\) | \(\ldots\) | \(F_{n-1}\) | \(F_n\) |
|---|---|---|---|---|---|---|---|---|---|---|---|
| \(y(k_1)\) | \(y(2)\) | \(\ldots\) | \(y(1)\) | \(\ldots\) | \(y(t)\) | \(\ldots\) | \(y(n-1)\) | \(y(n)\) |

We park cars \(F_1, \ldots, F_n\) using the preference sequence \(c\). The cars \(F_1, \ldots, F_{k_1-1}\) will take the spaces right after the trailer. The total length of \(F_1, \ldots, F_{k_1-1}\) is no more than the total length of \(C_1, \ldots, C_{k_1-1}\), and \(C_t\), since \(y(1) + y(2) + \cdots + y(k_1-1) + y(t) > y(2) + \cdots + y(k_1-1) + y(k_1)\). So \(F_{k_1}\) will park at the spot starting at \(c_k\), just as \(C_{k_1}\). But, as \(y(1) < y(k_1)\), after \(F_{k_1}\) is parked, the available space after \(F_{k_1}\) is nonempty and not a multiple of \(b\), while all the remaining cars are of length \(b\). Hence, it is not possible to park all of them and \(c \notin PS(\sigma_3(y_{inc}); z)\).

Combining Lemma 4.1 and Theorem 4.2, we obtain the following counting formula.

**Corollary 4.3.** Let \(z \in \mathbb{Z}_+\) and \(\vec{y} = (y_1, y_2, \ldots, y_n) \in \mathbb{Z}_+^n\). If \(\vec{y} \neq (s^n)\) for any integer \(s\), then

\[
\#SPS(\vec{y}; z) = z \cdot \prod_{i=1}^{n-1} (z + y(1) + \cdots + y(i)).
\]

where \(y(1) \leq y(2) \leq \cdots \leq y(n)\) is the order statistics of \(\vec{y}\).

## 5 Closing Remarks

In this paper, we studied increasing parking sequences and their connections with lattice paths. We also studied permutation-invariant parking sequences and length-invariant parking sequences. More precisely, we characterized the permutation-invariant parking sequences for some special families of length vectors. While it may not be easy to find a general formula for all cases, a natural direction to go would be to investigate other special cases of car lengths. Furthermore, in the study of parking functions we encounter quite a number of other mathematical structures including trees, non-crossing partitions, hyperplane arrangements, polytopes etc. It will be interesting to investigate if there is anything that connects other combinatorial structures to invariant parking sequences. Recently in [1], parking sequences were extended to the case in which one or more trailers are placed anywhere on the street alongside \(n\) cars with length vector \(\vec{y} = (1, 1, \ldots, 1)\). A natural generalization is to consider a similar scenario where \(\vec{y}\) is any length vector.
Acknowledgements

The authors are grateful for helpful conversations with Westin King which led to shorter proofs for Theorem 3.2 and Lemma 3.6. We also thank the two anonymous referees for carefully reading our manuscript and giving many helpful suggestions, which greatly improved the paper.

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