A labeled sequent calculus for propositional linear
time logic

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Abstract. A labeled sequent calculus \( \text{LSC} \) for propositional linear discrete time logic \( \text{PLTL} \) is introduced. Its sub-calculus \( \text{LSC}_{\text{TL}} \) is proved to be complete for some class of \( \text{PLTL} \) sequents.

Keywords: labeled sequent calculus, temporal logic.

1 Introduction

Temporal logic is a special type of modal logic. It provides a formal system for qualitatively describing and reasoning about how the truth values of assertions change over time. Propositional linear discrete time logic \( \text{PLTL} \) with temporal operators “next” and “always” is considered in the present paper.

Various syntactical proof-search systems are used for \( \text{PLTL} \). Some of them are:

- Sequent calculi with the invariant rule

\[
\frac{\Gamma \rightarrow \Delta, I; I \rightarrow \circ I; I \rightarrow A}{\Gamma \rightarrow \Delta, \Box A} (\rightarrow \Box_I),
\]

[10, 11]. There are some interesting works in which invariant-free (and cut-free) calculi for \( \text{PLTL} \) are constructed [3, 6].

- Sequent calculi with the infinitary rule

\[
\frac{\Gamma \rightarrow \Delta, A; \Gamma \rightarrow \Delta, \circ A; \ldots; \Gamma \rightarrow \Delta, \circ^n A; \ldots}{\Gamma \rightarrow \Delta, \Box A} (\rightarrow \Box_\omega),
\]

[12]. There are some interesting works concerning finitization of \( \omega \)-type rule \( (\rightarrow \Box_\omega) \) (see, e.g., [4]).

- Proof procedures containing loop-type axioms for logics sub-logic of which is propositional temporal branching time logic [9].

- Labeled sequent calculi [1, 2].

- Resolution-type proof procedures based on formulas in some normal form, see, e.g., [5].
In the present paper a labeled sequent calculus $\text{LSC}$ is presented. Its sub-calculus $\text{LSC}_{\text{TL}}$ is proved to be complete for some easily defined but large class of $\text{PLTL}$ sequents. Unlike the other deductive systems mentioned above, calculi $\text{LSC}$ and $\text{LSC}_{\text{TL}}$ are loop-axiom and invariant and infinitary rule free, which allows to construct effective proof-search procedures based on the calculi.

2 Syntax

Formulas are defined in the traditional way.

Formulas of the shape $x^k : A$, where $k \in \{0\} \cup \mathbb{N}$ (in particular, $x^0 = x$) and $A$ is a formula, are called labeled formulas, l-formulas for short; $x$ is called a label or a variable and $k$ its power. Labels/variables are denoted by $u$, $x$, $y$, $z$, $w$ and the corresponding powered labels by $u^k$, $x^k$, $y^k$, $z^k$, and $w^k$. The intended meaning of `$x : A$' is "$A$ holds at some moment of time $x$" and that one of `$x^k : A$' is "$A$ holds at the $k$-th from $x$ moment of time".

Expressions $m^n : A$, where $m, n \in \{0\} \cup \mathbb{N}$, are called fixed-label formulas and `$m^n$' fixed labels.

One more type of formulas is $x^m \leq y^n$, where $m \geq 0$. Such formulas are called order atoms.

Sequents are objects of the type $\Gamma \rightarrow \Delta$, where $\Gamma$ and $\Delta$ are some finite multisets of formulas.

Labeled sequents, l-sequents for short, are objects of the type $\Gamma \rightarrow \Delta$, where $\Gamma$ is some finite multiset of labeled formulas and order atoms; the same for $\Delta$ except that order atoms do not occur in it.

3 Semantics

Kripke semantics of $\text{PLTL}$ is defined as follows.

$([0] \cup \mathbb{N} \times \mathbf{P}) \to \{\top, \bot\}$, where $\mathbf{P}$ is the set of propositional variables.

$([0] \cup \mathbb{N} \times \mathbf{F}) \overset{\phi}{\to} \{\top, \bot\}$, where $\mathbf{F}$ is the set of formulas and $\phi$ is defined in the following way.

1. $\phi(i, E) = \tau(i, E)$, where $E$ is an atomic formula;
2. $\phi(i, A)$ is defined in the common way if $A$ is of the shape $\neg B$ or $B \theta C$, where $\theta$ is a logical connective;
3. $\phi(i, \circ A) = \top$ iff $\phi(i + 1, A) = \top$; otherwise, $\phi(i, \circ A) = \bot$;
4. $\phi(i, \Box A) = \top$ iff $\phi(j, A) = \top$ for all $j$ such that $j \geq i$; otherwise, $\phi(i, \Box A) = \bot$.

Some more notation:

(1) $(i^k : A) = \phi(i + k, A)$;
(2) $\models i^k : A$ iff $(i^k : A) = \top$ for any $\phi$;
(3) $\models x^k : A$ iff $\models i^k : A$ for all $i \geq 0$;
(4) $\models A$ iff $\phi(i, A) = \top$ for all $i \geq 0$ and every $\phi$

here $A$ is a label free formula, $k \geq 0$, and `$\models$' denotes validity.
A labeled sequent calculus for propositional linear time logic

Let \( \mathcal{L} \) be \( \{0\} \cup \mathbb{N} \), where \( \mathcal{L} \) is the set of labels.

The stable sequent \( S_\nu \) is obtained from \( S \) by substituting every label \( x_i \) by \( \nu(x_i) \). \( S_\nu \) is the set of labels and \( \varsigma \) is defined as follows:

if 

\[
S = x_1^1 : A_1, \ldots, x_k^k : A_k \rightarrow x_{k+1}^1 : B_{k+1}, \ldots, x_{k+m}^k : B_{k+m},
\]

then \( \varsigma(S_\nu) = \top \) if there are \( \phi, \nu, \) and \( t \), where \( 1 \leq t \leq (k + m) \), such that \( (\nu(x_i))^t : A_i = \perp \) or \( (\nu(x_i))^t : B_i = \top \).

A stable sequent \( S_\nu \) is valid, denoted by \( \models S_\nu \), if \( \varsigma(S_\nu) = \top \) for any \( \tau \).

A labeled sequent \( S_\nu \) is valid if it is of the shape \( \Gamma, \nu^k : E \rightarrow m^\nu : E, \Delta \), where \( l + k = m + n \).

A labeled sequent \( S \) is valid, denoted by \( \models S \), if every stable sequent obtained from \( S \) is valid.

4 Labeled sequent calculi LSC and LSC\(_{TL}^\sim\)

The labeled sequent calculus LSC for PLTL is defined as follows:

1. Axioms:

\[
\Gamma, x^k : E \rightarrow x^k : E, \Delta,
\]

where \( E \) is an atomic formula.

2. Logical rules:

\[
\begin{align*}
\frac{x^k : A, x^k : B ; \Gamma \rightarrow \Delta}{x^k : A \land B, \Gamma \rightarrow \Delta} (\land \rightarrow), & & \frac{\Gamma \rightarrow x^k : A, \Delta}{\Gamma \rightarrow x^k : A \land B, \Delta} (\rightarrow \land), \\
\frac{x^k : A \lor B, \Gamma \rightarrow \Delta}{x^k : A, \Gamma \rightarrow \Delta} (\lor \rightarrow), & & \frac{\Gamma \rightarrow x^k : A, x^k : B, \Delta}{\Gamma \rightarrow x^k : A \lor B, \Delta} (\rightarrow \lor), \\
\frac{\Gamma \rightarrow x^k : A, \Delta \quad x^k : A \rightarrow \neg A, \Gamma \rightarrow \Delta}{\Gamma \rightarrow x^k : A \lor B, \Delta} (\rightarrow \implies), & & \frac{\Gamma \rightarrow x^k : A, \Delta \quad \Gamma \rightarrow x^k : \neg A, \Delta}{\Gamma \rightarrow x^k : A, \Delta} (\rightarrow \neg), \\
\frac{\Gamma \rightarrow x^k : A, \Delta \quad \Gamma \rightarrow x^k : B, \Delta}{\Gamma \rightarrow x^k : A \lor B, \Delta} (\rightarrow \lor), & & \frac{\Gamma \rightarrow x^k : A, \Delta}{\Gamma \rightarrow x^k : A \lor B, \Delta} (\rightarrow \lor).
\end{align*}
\]

Here \( A \) and \( B \) arbitrary formulas.

3. Temporal rules:

\[
\begin{align*}
\frac{\Gamma \rightarrow x^{k+1} : A, \Delta}{\Gamma \rightarrow x^k : \Diamond A, \Delta} (\rightarrow \Diamond), & & \frac{x^{k+1} : A, \Gamma \rightarrow \Delta}{x^k : \Diamond A, \Gamma \rightarrow \Delta} (\rightarrow \Diamond), \\
\frac{x \leq y, \Gamma \rightarrow y^k : A, \Delta}{\Gamma \rightarrow x^k : \Diamond A, \Delta} (\rightarrow \Diamond), & & \frac{y^{k+m} : A, x^{k+m} \leq y^{k+m}, x^k : \Diamond A, \Gamma \rightarrow \Delta}{x^{k+m} \leq y^{k+m}, x^k : \Diamond A, \Gamma \rightarrow \Delta} (\rightarrow \Diamond).
\end{align*}
\]

Here \( k, m \geq 0 \); \( y \) in \( (\rightarrow \Diamond) \) does not occur in the conclusion.
4. Rules for order atoms:

\[
\frac{x \leq y, \Gamma \rightarrow \Delta}{\Gamma \rightarrow \Delta} \quad \text{Ref}, \quad \frac{x^{k+1} \leq y^{k+1}, x^{k} \leq y^{k}, \Gamma \rightarrow \Delta}{\Gamma \rightarrow \Delta} \quad \text{Fwd}^{+1},
\]

\[
\frac{x \leq z, x \leq y, y \leq z, \Gamma \rightarrow \Delta}{\Gamma \rightarrow \Delta} \quad \text{Trans},
\]

\[
\frac{y \leq z, x \leq y, x \leq z, \Gamma \rightarrow \Delta}{x \leq y, z \leq x, \Gamma \rightarrow \Delta} \quad \text{Lin}.
\]

Here \(x, y, z\) are unequal in pairs in \(\text{Trans}\) and \(\text{Lin}\); \(x \leq x\) does not occur in \(\Gamma\) in \(\text{Ref}\); \(x^{k+1} \leq y^{k+1}\) does not occur in \(\Gamma\) in \(\text{Fwd}\); \(x \leq z\) does not occur in \(\Gamma\) in \(\text{Trans}\); \(\text{In Lin}, \) neither \(y \leq z\) nor \(z \leq y\) occur in \(\Gamma\) neither can be obtained by some backward applications of \(\text{Trans}\).

The calculus \(\text{LSC}_{\text{TL}}\) is obtained from \(\text{LSC}\) by dropping \(\text{Trans}\) and \(\text{Lin}\).

A formula \(F\) is called derivable in the labeled sequent calculus \(\text{LSC}(\text{LSC}_{\text{TL}})\) iff \(\text{LSC}(\text{LSC}_{\text{TL}}) \vdash x : F\).

A sequent \(S\) is called derivable in \(\text{LSC}(\text{LSC}_{\text{TL}})\) iff \(\text{LSC}(\text{LSC}_{\text{TL}}) \vdash x : S\).

The Hilbert-style calculus \(\text{HSC}\) for \(\text{PLTL}\) is defined by axioms:

\(A_0\): propositional tautologies; \(A_1\): \(\neg p \equiv \neg \neg p\);

\(A_2\): \(\neg p \equiv (\neg p \cap q) ; \quad A_3\): \(\Box (p \cap q) \equiv (\Box p \cap \Box q)\);

\(A_4\): \(\Box p \equiv p ; \quad A_5\): \(\Box (p \cap q) \equiv (\Box p \cap \Box q)\);

\(A_6\): \(p \cap q \equiv q \cap p\).

and derivation rules:

\[
\frac{p}{\Box p}, \quad \frac{p}{\Box p}, \quad \frac{p \cap q}{q \cap p},
\]

where \(p\) and \(q\) are arbitrary \(\text{PLTL}\) formulas. It is well known that this calculus is sound and complete for \(\text{PLTL}\), see, e.g. [7].

5 Some Properties of LSC and \(\text{LSC}_{\text{TL}}\)

**Lemma 1.** If \(\text{LSC}(\text{LSC}_{\text{TL}}) \vdash V S\), then \(\text{LSC}(\text{LSC}_{\text{TL}}) \vdash V' S(w/u)\) and \(h(V') \leq h(V)\), where \(S(w/u)\) is obtained from \(S\) by substituting the label \(w\) for the label \(u\).

A rule is height-preserving admissible if, whenever its premiss(es) is (are) derivable, also its conclusion is derivable with the same bound on the derivation height.

**Lemma 2.** The rule of weakening

\[
\frac{\Gamma \rightarrow \Delta}{\Gamma', \Gamma \rightarrow \Delta, \Delta'}\quad (w)
\]

is height-preserving admissible in \(\text{LSC}\) and \(\text{LSC}_{\text{TL}}\).

A rule is height-preserving invertible if, whenever its conclusion is derivable, also its premiss(es) is (are) derivable with the same bound on the derivation height.

**Lemma 3.** All \(\text{LSC}\) rules are height-preserving invertible in \(\text{LSC}\), and all \(\text{LSC}_{\text{TL}}\) rules are height-preserving invertible in \(\text{LSC}_{\text{TL}}\).
A labeled sequent calculus for propositional linear time logic

Lemma 4. The rules of contraction

\[ \frac{C, C, \Gamma \rightarrow \Delta}{C, \Gamma \rightarrow \Delta} \quad \text{(c \rightarrow)} \]
\[ \frac{\Gamma \rightarrow \Delta, C, C}{\Gamma \rightarrow \Delta, C} \quad \text{(\rightarrow c)} \]

are height-preserving admissible in LSC and LSC\(_{\text{TL}}\).

A sequent \( S \) is called proper if the fact that \( x^k \leq y^k \) occurs in \( S \), where \( k \geq 0 \), implies that \( x^0 \leq y^0 \) occurs in \( S \).

Theorem 1. The rule of cut

\[ \frac{\Pi \rightarrow C, \Lambda; C, \Gamma \rightarrow \Delta \quad \Pi, \Gamma \rightarrow \Lambda, \Delta}{\Pi, \Gamma \rightarrow \Lambda, \Delta} \quad \text{cut} \]

is admissible in LSC and LSC\(_{\text{TL}}\), where the premises are proper.

Lemma 5. All LSC rules are correct: if the premise(s) is (are) valid, then so is the conclusion. In addition, if the conclusion is valid, then so is (are) the premise(s).

Lemma 6. Any labeled sequent of the shape \( \Gamma, x : A \rightarrow x : A, \Delta \) is derivable in LSC and LSC\(_{\text{TL}}\).

Theorem 2. HSC \( \vdash d \) \( F \) implies \( LSC_{\text{TL}} \vdash \rightarrow x : F \), where the induction axiom \( A_6 \) is not used in \( d \).

If \( \Gamma = A_1, \ldots, A_n \), then \( \theta \Gamma = (A_1 \theta \ldots \theta A_n) \), where \( \theta \in \{\land, \lor\} \). If \( S = \Gamma \rightarrow \Delta \), then \( F(S) = \neg(\land \Gamma) \lor (\lor \Delta) \).

By Theorem 2 and invertibility of the rules \( \rightarrow \lor \), \( \rightarrow \), and \( \land \rightarrow \), LSC\(_{\text{TL}}\) is complete for sequents \( S = \Gamma \rightarrow \Delta \) such that \( F(S) \) is derivable in HSC without using the axiom \( A_6 \).

An example of non-derivable in LSC\(_{\text{TL}}\) formula is

\[ \square A \supset \square \square A. \]

Theorem 2 implies that this formula is not derivable in HSC without the axiom \( A_6 \).

This formula is derivable in LSC.

Some examples of non-derivable in LSC formulas are

\[ (A \land \square A) \supset \square A \quad \text{and} \quad (A \land (A \supset \square A)) \supset \square A. \]

We get by Theorem 2 that these formulas are not derivable in HSC without \( A_6 \).

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R. Alonderis

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REZIUMĖ

Žymėtas sekvencinis skaičiavimas propozicinei tiesinio laiko logikai

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Darbe yra pateiktas žymėtas sekvencinis skaičiavimas propozicinei tiesinio laiko logikai. Įrodyta, kad šis skaičiavimas yra pilnas tam tikros nagrinėjamos logikos sekvencijų klasės atžvilgiu.

Raktiniai žodžiai: žymėtas sekvencinis skaičiavimas, laiko logika.