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A NOTE ON THE STABILITY NUMBER OF AN ORTHOGONALITY GRAPH

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A note on the stability number of an orthogonality graph

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Abstract

We consider the orthogonality graph $\Omega(n)$ with $2^n$ vertices corresponding to the vectors $\{0, 1\}^n$, two vertices adjacent if and only if the Hamming distance between them is $n/2$. We show that, for $n = 16$, the stability number of $\Omega(n)$ is $\alpha(\Omega(16)) = 2304$, thus proving a conjecture by Galliard [7]. The main tool we employ is a recent semidefinite programming relaxation for minimal distance binary codes due to Schrijver [16].

Moreover, we give a general condition for Delsarte bound on the (co)cliques in graphs of relations of association schemes to coincide with the ratio bound, and use it to show that for $\Omega(n)$ the latter two bounds are equal to $2^n/n$.

Keywords: Semidefinite programming, minimal distance codes, stability number, orthogonality graph, Hamming association scheme, Delsarte bound.

AMS subject classification: 90C22, 90C27, 05C69,05C15,

JEL code: C0, C61

1 Introduction

The graph $\Omega(n)$ and its properties

Let $\Omega(n)$ be the graph on $2^n$ vertices corresponding to the vectors $\{0, 1\}^n$, such that two vertices are adjacent if and only if the Hamming distance between them is $n/2$. Note that $\Omega(n)$ is $k$-regular, where $k = \binom{n}{\frac{n}{2}}$.

It is known that $\Omega(n)$ is bipartite if $n = 2 \mod 4$, and empty if $n$ is odd. We will therefore assume throughout that $n$ is a factor of 4. The graph owns its

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name to another description, in terms of ±1-vectors. Then the orthogonality of vectors corresponds to the Hamming distance $n/2$.

Moreover, $\Omega(n)$ consists of two isomorphic components, one containing all the vertices of even Hamming weight and the other the vertices of odd Hamming weight. For a detailed discussion of the properties of $\Omega(n)$, see Godsil [9] and the PhD thesis of Newman [13].

In this note we study upper bounds on the stability number $\alpha(\Omega(n))$. Galliard [7] pointed out the following way of constructing maximal stable sets in $\Omega(n)$. Consider the "odd component" of $\Omega(n)$ and take all vertices of Hamming weight 1, 3, $\ldots$, $n/4 - 1$. Obviously, these vertices form a stable set of $\Omega(n)$ of size

$$\sum_{1 \leq i \leq n/8} \binom{n}{2i - 1}. \quad (1)$$

We can double the size of this stable set by adding the bit-wise complements of the vertices in $S$, and double it again by taking the union with the corresponding stable set in the other (isomorphic) component.

Thus we find that

$$\alpha(\Omega(n)) \geq 4 \sum_{1 \leq i \leq n/8} \binom{n}{2i - 1} := \alpha(n).$$

For $n = 16$ this evaluates to $\alpha(\Omega(n)) \geq 2304$. Galliard et al [8] could show that $\alpha(\Omega(16)) \leq 3912$. In this note we will show that, in fact, $\alpha(\Omega(16)) = 2304$. This
was conjectured by Galliard [7], and Newman [13] has recently conjectured that the value (1) actually equals $\alpha(\Omega(n))$ whenever $n$ is a multiple of 4.

**A quantum information game**

One motivation for studying the graph $\Omega(n)$ comes from quantum information theory. Consider the following game from [8].

Let $n \geq 1$ and $N = 2^n$. Two players, A and B, are asked the questions $x_A$ and $x_B$, coded as $N$-bit strings satisfying

$$d_H(x_A, x_B) \in \left\{0, \frac{1}{2} N\right\}$$

where $d_H$ denotes the Hamming distance. A and B win the game if they give answers $y_A$ and $y_B$, coded as binary strings of length $n$ such that

$$y_A = y_B \iff x_A = x_B.$$  

A and B are not allowed any communication (except a priori deliberation). It is known that A and B can always win the game if their $n$ output bits are *maximally entangled quantum bits* [2] (see also [13]).

For classical bits, it was shown by Galliard et al [8] that the game cannot always be won if $n = 4$. The authors proved this by pointing out that whether or not the game can always be won is equivalent to the question

$$\chi(\Omega(n)) \leq N \equiv 2^n?$$

Indeed, if $\chi(\Omega(n)) \leq N$ then A and B may color $\Omega(n)$ a priori using $N$ colors. The questions $x_A$ and $x_B$ may then be viewed as two vertices of $\Omega(n)$, and the A and B may answer their respective questions by giving the color of the vertices $x_A$ and $x_B$ respectively, coded as binary strings of length $\log_2 N = n$.

Galliard et al. [8] showed that $\chi(\Omega(16)) > 16$, i.e. that the game cannot be won for $n = 16$. They proved this by showing that $\alpha(\Omega(16)) \leq 3912$ which implies

$$\chi(\Omega(16)) \geq \left\lceil \frac{2^{16}}{\alpha(\Omega(16))} \right\rceil \geq \left\lceil \frac{2^{16}}{3912} \right\rceil = 17.$$  

In this note we sharpen their bound by showing that $\alpha(\Omega(16)) = 2304$, which implies $\chi(\Omega(16)) \geq 21$.

Our main tool will be a semidefinite programming bound on $\alpha(\Omega(n))$ that is due to Schrijver [16], where it is formulated for minimal distance binary codes.

**2 Upper bounds on $\alpha(\Omega(n))$**

In this section we give a review of known upper bounds on $\alpha(\Omega(n))$ and their relationship.
2.1 The ratio bound

The following discussion is condensed from Godsil [9].

**Theorem 1.** Let $G = (V, E)$ be a $k$-regular graph with adjacency matrix $A(G)$, and let $\lambda_{\text{min}}(A(G))$ denote the smallest eigenvalue of $A(G)$. Then

$$\alpha(G) \leq \frac{|V|}{1 - \frac{k}{\lambda_{\text{min}}(A(G))}}.$$  \hspace{2cm} (2)

This bound is called the *ratio bound*, and was first derived by Delsarte [4] for graphs in association schemes (see Sect. 2.2 for more on the latter). Recall that $\Omega(n)$ is $k$-regular with $k = \binom{n}{\frac{1}{2}n}$. Ignoring multiplicities, the spectrum of $\Omega(n)$ is given by

$$\lambda_m = \frac{2^m}{\left(\frac{1}{2}n\right)!}(m-1)(m-3)\cdots(m-n+1) \quad (m = 1, \ldots, n).$$  \hspace{2cm} (3)

The minimum is reached at $m = 2$, and we get

$$\lambda_{\text{min}}(A(\Omega(n))) = \frac{2^2}{\left(\frac{1}{2}n\right)!}(1)(-1)(-3)\cdots(-n+3) = -\frac{n}{n-1}. \quad (4)$$

The ratio bound therefore becomes

$$\alpha(\Omega(n)) \leq \frac{2^n}{n}. \quad (5)$$

This is the best known upper bound on $\alpha(\Omega(n))$, but it is known that this bound is not tight: Frankl and Rödl [6] showed that there exists some $\epsilon > 0$ such that $\alpha(\Omega(n)) \leq (2 - \epsilon)^n$. For specific (small) values of $n$ one can improve on the bound (5), as we will show for $n \leq 32$.

2.2 The Delsarte bound and $\vartheta'$

Here we are going to use more linear algebra that naturally arise around $\Omega(n)$. We recall the following definitions, cf. e.g. Bannai and Ito [1].

**Association schemes.** An association scheme $\mathcal{A}$ is a commutative subalgebra of the full $v \times v$-matrix algebra with a distinguished basis ($A_0 = I, A_1, \ldots, A_n$) of 0-1 matrices. One often views $A_j$, $j \geq 1$, as the adjacency matrix of a graph on $v$ vertices; $A_j$ is often referred to as the $j$-th *relation* of $\mathcal{A}$. As the $A_j$’s commute, they have $n + 1$ common eigenspaces $V_i$. Then $\mathcal{A}$ is isomorphic, as an algebra, to the algebra of diagonal matrices $\text{diag} (P_{0,j}, \ldots, P_{n,j})$, where $P_{ij}$ denotes the eigenvalue of $A_j$ on $V_i$. The matrix $P = (P_{ij})$ is called *first eigenvalue matrix* of $\mathcal{A}$. The set of $A_j$’s is closed under taking transpositions: for each $0 \leq j \leq n$ there exists $j'$ so that $A_j = A_{j'}^T$. In particular, $P_{ij} = P_{ij'}$. An association scheme with all $A_j$ symmetric is called *symmetric*, and here we shall consider such
schemes only. There is a matrix \( Q \) (called second eigenvalue matrix) satisfying \( PQ = QP = vI \). In what follows it is assumed (as is customary in the literature) that the eigenspace \( V_0 \) corresponds to the eigenvector \((1, \ldots, 1)\); then the 0-th row of \( P \) consists of the degrees \( v_j \) of the graphs \( A_j \). It is remarkable that the 0-th row of \( Q \) consists of dimensions of \( V_i \).

Let \( \vartheta' \) denote the Schrijver \( \vartheta' \)-function [15]:

\[
\vartheta'(G) = \max \{ \text{Tr} (JX) : \text{Tr} (AX) = 0, \text{Tr}(X) = 1, X \geq 0, X \geq 0 \}.
\]

For any graph \( G \) one has \( \alpha(G) \leq \vartheta'(G) \). Moreover, \( \vartheta'(G) \) is smaller than or equal to the ratio bound (2) for regular graphs, as noted by Godsil [9, Sect. 3.7].

For graphs with adjacency matrices of the form \( \sum_{j \in M} A_j \), with \( M \subset \{1, \ldots, n\} \) and \( A_j \)'s from the 0-1 basis of an association scheme \( A \), the bound \( \vartheta' \) coincides, as was proved by Schrijver [15], with the following bound due to Delsarte [3, 4]

\[
\max 1^T w \text{ subject to } w \geq 0, Q^T w \geq 0, w_0 = 1, w_j = 0 \text{ for } j \in M,
\]  

(6)

where \( Q \) is the second eigenvalue matrix of \( A \).

The bound (6) is often stated for (and was originally developed for) bounding the maximal size of a \( q \)-ary code of length \( n \) and minimal distance \( d \); then the association scheme \( A \) becomes the Hamming distance association scheme \( H(n, q) \) and \( M = \{1, \ldots, d - 1\} \). The relations of \( H(n, q) \) can be viewed as graphs on the vertex set of \( n \)-strings on \( \{0, \ldots, q - 1\} \): the \( j \)-th graph of \( H(n, q) \) is given by

\[
(A_j)_{XY} = \begin{cases} 
1 & \text{if } d_H(X, Y) = j \\
0 & \text{otherwise}
\end{cases}
\]

For \( H(n, q) \) the first and the second eigenvalue matrices \( P \) and \( Q \) coincide, and are given by \( P_{ij} = K_i(j) \), where \( K_k \) is the Krawtchouk polynomial

\[
K_k(x) := \sum_{j=0}^{n} (-1)^j (q - 1)^{k-j} \binom{x}{j} \binom{n-x}{k-j}.
\]

For \( \Omega(n) \), the bound (6) is as above with \( A = H(n, 2) \) and \( M = \{\frac{n}{2}\} \). Newman [13] has shown computationally that \( \vartheta'(\Omega(n)) = 2^n/n \) if \( n \leq 64 \), i.e. the ratio and \( \vartheta' \) bounds coincide for \( \Omega(n) \) if \( n \leq 64 \). We show that it is the case for all \( n \), as an easy consequence of the following.

**Proposition 1.** Let \( A \) be an association scheme with the 0-1 basis \( (A_0, \ldots, A_n) \) and eigenvalue matrices \( P \) and \( Q \). Let \( A_r \) have the least eigenvalue \( \tau = P_{rr} \) and assume

\[
v_r P_{ri} \geq v_i \tau, \quad 0 \leq i \leq n.
\]

Then the Delsarte bound (6), with \( M = \{r\} \), and the ratio bound (2) for \( A_r \) coincide.

**Proof.** Let \( P_j \) denote the \( j \)-th row of \( P \).
As we already mentioned, the bound (2) for regular graphs always majorates (6). Thus it suffices to present a feasible vector for the LP in (6) that gives the objective value the same as (2).

We claim that
\[ a = \frac{-\tau}{v_r - \tau} P_0 + \frac{v_r}{v_r - \tau} P_\ell \]
is such a vector. It is straightforward to check that \( a_0 = 1 \) and \( a_r = 0 \), as required. By the assumption of the proposition, \( a \geq 0 \). As \( PQ = vI \), any nonnegative linear combination \( z \) of the rows of \( P \) satisfies \( Q^T z \geq 0 \). As \( a \) is such a combination, we obtain \( Q^T a \geq 0 \).

Finally, to compute \( 1^T a \), note that \( 1^T P_0 = v \) and \( 1^T P_\ell = 0 \).

**Corollary 1.** The bounds (6) and (2) coincide for \( \Omega(n) \).

**Proof.** We apply Proposition 1 to \( A = H(n, 2) \) and \( r = \frac{n}{2} \). Then the eigenvalues of \( A_r = \Omega(n) \) given in (3) comprise the \( r \)-th column on \( P \), in particular the least eigenvalue \( \tau \) equals \( P_{2,r} \), by (4) above. The assumption of the proposition translates into
\[
\left( \begin{array}{c}
\binom{n}{2} \\
\frac{n}{2}
\end{array} \right) K_i(2) - \left( \begin{array}{c}
\binom{n}{i} \\
\frac{n}{2}
\end{array} \right) K_{\frac{n}{2}}(2) = \frac{2^{\frac{n}{2}+2}(n-2)!(n-1)!(\frac{n}{2} - i)^2}{d!(\frac{n}{2})!(n-i)!} \geq 0,
\]
as claimed.

\[ \square \]

### 2.3 Schrijver’s improved SDP-based bound

Recently, Schrijver [16] has suggested a new SDP-based bound for minimal distance codes, that is at least as good as the \( \vartheta' \) bound, and still of size polynomial in \( n \). It is given as the optimal value of a semidefinite programming (SDP) problem.

In order to introduce this bound (as applied to \( \alpha(\Omega(n)) \)) we require some notation.

For \( i, j, t \in \{0, 1, \ldots, n\} \), and \( X, Y \in \{0, 1\}^n \) define the matrices
\[
(M^{t}_{i,j})_{X,Y} = \begin{cases} 
1 & \text{if } |X| = i, |Y| = j, d_H(X,Y) = n-t \\
0 & \text{otherwise}
\end{cases}
\]

The upper bound is given as the optimal value of the following semidefinite program:

\[
\bar{\alpha}(n) := \max \sum \binom{n}{i} x^0_{i,0}
\]

subject to
\[
\begin{align*}
x^0_{0,0} &= 1 \\
0 &\leq x^t_{i,j} \leq x^0_{i,0} \text{ for all } i, j, t \in \{0,\ldots,n\} \\
x^t_{i,j} &= x^t_{i',j'} \text{ if } \{i', j', i' + j' - 2t'\} \text{ is a permutation of } \{i, j, i + j - 2t\} \\
x^t_{i,j} &= 0 \text{ if } \{i, j, i + j - 2t\} \cap \{\frac{1}{2}n\} \neq \emptyset,
\end{align*}
\]

\[1\text{Here } m!! = m(m-2)(m-4)\ldots, \text{ the double factorial}.\]
as well as
\[
\sum_{i,j,t} x^t_{i,j} M^t_{i,j} \geq 0, \quad \sum_{i,j,t} (x^0_{i+j-2t,0} - x^t_{i,j}) M^t_{i,j} \geq 0.
\]

The matrices \(M^t_{i,j}\) are of order \(2^n\) and therefore too large to compute with in general. Schrijver pointed out that these matrices form a basis of the Terwilliger algebra of the Hamming scheme, and worked out the details for computing the irreducible block diagonalization of this (non-commutative) matrix algebra of dimension \(O(n^3)\).

Thus, analogously to the \(\vartheta'-\)case, the constraint \(\sum_{i,j,t} x^t_{i,j} M^t_{i,j} \geq 0\) is replaced by
\[
\sum_{i,j,t} x^t_{i,j} Q^T M^t_{i,j} Q \geq 0
\]
where \(Q\) is an orthogonal matrix that gives the irreducible block diagonalization.

For details the reader is referred to Schrijver [16]. Since SDP solvers can exploit block diagonal structure, this reduces the sizes of the matrices in question to the extent that computation is possible in the range \(n \leq 32\).

### 2.4 Laurent’s improvement

In Laurent [12] one finds a study placing the relaxation [16] into the framework of moment sequences of [10, 11]. This study also explains the relationship with known lift-and-project methods for obtaining hierarchies of upper bounds on \(\alpha(G)\).

Moreover, Laurent [12] suggests a refinement of the Schrijver relaxation that takes the following form:
\[
l_+(n) := \max 2^n x^0_{0,0}
\]
subject to
\[
0 \leq x^t_{i,j} \leq x^0_{i,0} \text{ for all } i, j, t \in \{0, \ldots, n\}
\]
\[
x^t_{i,j} = x^t_{i',j'} \text{ if } \{i', j', i' + j' - 2t'\} \text{ is a permutation of } \{i, j, i+j-2t\}
\]
\[
x^t_{i,j} = 0 \text{ if } \{i, j, i+j-2t\} \cap \{\frac{1}{2}n\} \neq \emptyset,
\]

as well as
\[
\sum_{i,j,t} x^t_{i,j} M^t_{i,j} \geq 0
\]
and
\[
\left(1 - x^0_{0,0} \frac{c^T}{c} \sum_{i,j,t} (x^0_{i+j-2t,0} - x^t_{i,j}) M^t_{i,j}\right) \geq 0,
\]
where \(c := \sum_{t=0}^n (x^0_{0,0} - x^0_{0,t}) \chi_t\), and \(\chi_t\) is defined by
\[
(\chi_t)X := \begin{cases} 1 & \text{ if } |X| = i \\ 0 & \text{ else.} \end{cases}
\]

This SDP problem may be block-diagonalised as before to obtain an SDP of size \(O(n^3)\).
3 Computational results

To summarize, the bounds we have mentioned satisfy:

$$\alpha(n) \leq \alpha(\Omega(n)) \leq l^+(n) \leq \bar{\alpha}(n) \leq \vartheta'(\Omega(n)) = 2^n/n.$$ 

In Table 1 we show the numerical values for $\bar{\alpha}(n)$ and $l^+(n)$ that were obtained using the SDP solver SeDuMi by Sturm [17], with Matlab 7 on a Pentium IV machine with 1GB of memory. The Matlab routines that we have written to generate the corresponding SeDuMi input are available online [14].

| $n$ | $\alpha(n)$ | $l^+(n)$ | $\bar{\alpha}(n)$ | $\vartheta'(\Omega(n)) = \lfloor 2^n/n \rfloor$ |
|-----|-------------|-----------|-------------------|---------------------------------|
| 16  | 2304        | 2304      | 2304              | 4096                            |
| 20  | 20,144      | 20,166.62 | 20,166.98         | 52,428                          |
| 24  | 178,208     | 183,373   | 184,194           | 699,050                         |
| 28  | 406,336     | 1,883,009 | 1,848,580         | 9,586,980                       |
| 32  | 14,288,896  | 21,103,609| 21,723,404        | 134,217,728                     |

Table 1: Lower and upper bounds on $\alpha(\Omega(n))$.

Note that the lower and upper bounds coincide for $n = 16$, proving that $\alpha(\Omega(16)) = 2304$. The best previously known upper bound, obtained by an ad hoc method, was $\alpha(\Omega(16)) \leq 3912$ [8].

The value $\bar{\alpha}(20) = 20,166.98$ implies that

$$\alpha(\Omega(20)) \in \{20144, 20148, 20152, 20156, 20160, 20164\}$$

since $\alpha(\Omega(n))$ is always a factor of 4. Another implication is that $n = 20$ is the smallest value of $n$ where the upper bounds $\bar{\alpha}(n)$ and $l^+(n)$ are not tight. It is worth noticing that the Schrijver and Laurent bounds ($\bar{\alpha}(n)$ and $l^+(n)$ respectively) give relatively big improvements over the Delsarte bound $\frac{2^n}{n}$. This is in contrast to the relatively small improvements that these bounds give for binary codes, cf. [16, 12]. We also note that these relaxations are numerically ill-conditioned for $n \geq 24$. This makes it difficult to solve the corresponding SDP problems to high accuracy. The recent study by De Klerk, Pasechnik, and Schrijver [5] suggests a different way to solve such SDP problems, leading to larger SDP instances, but which may avoid the numerical ill-conditioning caused by performing the irreducible block factorization.

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