ON THE ENTROPY OF RANDOM FIBONACCI WORDS

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Abstract

The random Fibonacci chain is a generalisation of the classical Fibonacci substitution and is defined as the rule mapping \(0 \mapsto 1\) and \(1 \mapsto 01\) with probability \(p\) and \(1 \mapsto 10\) with probability \(1 - p\) for \(0 < p < 1\) and where the random rule is applied each time it acts on a 1. We show that the topological entropy of this object is given by the growth rate of the set of inflated random Fibonacci words.

MSC2010 classification: 68R15 Combinatorics on words, 05A16 Asymptotic enumeration, 37B10 Symbolic dynamics.

1 Introduction

In [2] Godrèche and Luck define the random Fibonacci chain by the generalised substitution

\[ \theta : \begin{cases} 
0 & \mapsto 1 \\
1 & \mapsto \{ \begin{align*} 
& 01 \text{ with probability } p \\
& 10 \text{ with probability } 1 - p
\end{align*} \} \end{cases} \]  

for \(0 < p < 1\) and where the random rule is applied each time \(\theta\) acts on a 1. They introduce the random Fibonacci chain when studying quasi-crystalline structures and tilings in the plane. In their paper it is claimed without proof that the topological entropy of the random Fibonacci chain is given by the growth rate of the set of inflated random Fibonacci words. We give here a proof of this fact.

Before we can state our main theorem we need to introduce some notation. A word \(w\) over an alphabet \(\Sigma\) is a finite sequence \(w_1w_2\ldots w_n\) of symbols from \(\Sigma\). We let here \(\Sigma = \{0,1\}\). We denote a sub-word or a factor of \(w\) by \(w[a,b] = w_aw_{a+1}w_{a+2}\ldots w_{b-1}w_{b}\) and similarly we let \(W[a,b] = \{w[a,b] : w \in W\}\). By \(|\cdot|\) we mean the length of a word and the cardinality of a set. Note that \(|w[a,b]| = b - a + 1\).

For two words \(u = u_1u_2u_3\ldots u_n\) and \(v = v_1v_2v_3\ldots v_m\) we denote by \(uv\) the concatenation of the two words, that is, \(uv = u_1u_2u_3\ldots u_nv_1v_2\ldots v_m\). Similarly we let for two sets of words \(U\) and \(V\) their product be the set \(UV = \{uv : u \in U, v \in V\}\) containing all possible concatenations.

Letting \(\theta\) act on the word 0 repeatedly yields us an infinite sequence of words \(r_n = \theta^{n-1}(0)\). We know that \(r_1 = 0\) and \(r_2 = 1\). But \(r_3\) is either 01 or
10, with probability $p$ or $1 - p$ respectively. The sequence $\{r_n\}_{n=1}^{\infty}$ converges in distribution to an infinite random word $r$. We say that $r_n$ is an inflated word (under $\theta$) in generation $n$ and we introduce here sets that correspond to all inflated words in generation $n$;

**Definition 1.** Let $A_0 := \emptyset$, $A_1 := \{0\}$ and $A_2 := \{1\}$ and for $n \geq 3$ we define recursively

$$A_n := A_{n-1}A_{n-2} \cup A_{n-2}A_{n-1}$$

and we let $A := \lim_{n \to \infty} A_n$.

We shall later on, as a direct consequence of Proposition 7, see that the set $A$ is a well defined set. It is clear from the definition of $A_n$ that all elements in $A_n$ have the same length, that is, for all $x, y \in A_n$ we have $|x| = |y|$. We shall frequently use the length of the elements in $A_n$, and therefore we introduce the notation

$$f_n := |x| \text{ for } x \in A_n, \ n \geq 1$$

and for completeness we set $f_0 := 0$. From the recursion (2) and the definition of $f_n$ we have immediately the following proposition

**Proposition 2.** The numbers $f_n$ are the Fibonacci-numbers, that is, $f_0 = 0$, $f_1 = 1$ and $f_n = f_{n-1} + f_{n-2}$ for $n \geq 2$.

For a word $w$ we say that $x$ is a factor of $w$ if there are two words $u, v$ such that $w = uxv$. The factor set $F(S, n)$, of a set of words $S$, is the set of all factors of length $n$ of the words in $S$. We introduce the abbreviated notation $F_n$ for the special factor sets of the random Fibonacci words,

$$F_n := F(A, f_n).$$

It seems to be a hard problem to give an explicit expression for the size of $|F_n|$. We have to leave that question open, but we shall give a rough upper bound of $|F_n|$.

The topological entropy of the random Fibonacci chain is defined as the limit $\lim_{n \to \infty} \frac{1}{n} \log |F(A, n)|$. The existence of this limit is direct by Fekete’s lemma [1] since we have sub-additivity, $\log |F(A, n + m)| \leq \log |F(A, n)| + \log |F(A, m)|$. We can now state the main result in this paper

**Theorem 3.** The logarithm of the growth rate of the size of the set of inflated random Fibonacci words equals the topological entropy of the random Fibonacci chain, that is

$$\lim_{n \to \infty} \frac{\log |A_n|}{f_n} = \lim_{n \to \infty} \frac{\log |F_n|}{f_n}.$$  

The outline of the paper is that we start by studying the set $A_n$ and restate some of the results obtained by Godrèche and Luck. Thereafter we continue by looking at sets of prefixes of the set $A_n$. Next we give a finite method for finding the factor set $F_n$ and finally we present an estimate of $|F_n|$ in terms of $|A_n|$, leading to the proof of Theorem 3.
In this section we present the sets of inflated random Fibonacci words and give an insight to their structure. For some small values of \( n \) we have directly from the definition

\[
A_1 = \{0\}, \quad A_2 = \{1\}, \quad A_3 = \{01, 10\}, \quad A_4 = \{011, 101, 110\}, \\
A_5 = \{01011, 01101, 01110, 10011, 10101, 10110, 11001, 11010\}.
\]

For a word \( x \) we denote the reversed word of \( x \) by \( x^r \), that is, if \( x = x_1x_2 \ldots x_n \) then \( x^r = x_nx_{n-1} \ldots x_1 \). Similarly we write \( A_n^r = \{x^r : x \in A_n\} \).

**Proposition 4.** For \( n \geq 1 \) we have \( A_n = A_n^r \).

**Proof.** We give a proof by induction on \( n \). The basis cases \( n = 1 \) and \( n = 2 \) are clear as \( A_1 = \{0\} \) and \( A_2 = \{1\} \). Now assume for induction that the statement of the proposition holds for \( 2 \leq n \leq p \). Using (2) and the induction assumption give

\[
A_p^r + A_{p-1}^r = (A_p A_{p-1} \cup A_{p-1} A_p)^r \\
= A_p^r A_{p-1}^r \cup A_{p-1}^r A_p^r \\
= A_{p-1} A_p \cup A_p A_{p-1} \\
= A_{p+1},
\]

which completes the induction.

Now let us turn to the question of the size of \( A_n \). In the recursive definition (2) of \( A_n \) we have that the overlap is

\[
A_n A_{n-1} \cap A_{n-1} A_n = A_{n-2} A_{n-3} A_{n-2}.
\]

Knowing this we obtain the recursion for \( n \geq 3 \)

\[
|A_n| = 2|A_{n-1}||A_{n-2}| - |A_{n-2}|^2|A_{n-3}|,
\]

with \( |A_0| = 0 \), \( |A_1| = 1 \) and \( |A_2| = 1 \). The recursion (4) was given by Godrèche and Luck in [2]. They also gave the following proposition, which we for completeness restate here with its short proof.

**Proposition 5 ([2]).** With \( |A_1| = 1 \) and \( |A_2| = 1 \) we have for \( n \geq 3 \)

\[
|A_n| = \frac{n-1}{n-2}|A_{n-1}| |A_{n-2}|.
\]
Proof. Let \( a_n = |A_n|/(|A_{n-1}|A_{n-2}) \). Then we can rewrite the recursion (4) into
\[
a_n = 2 - \frac{1}{a_{n-1}}
\]
with the initial condition \( a_3 = 2 \). We claim that \( a_n = \frac{n-1}{n-2} \). The proof of the claim is a straightforward induction. The basis is clear, and if we assume the claim being true for \( n = p \) we get for \( n = p + 1 \)
\[
a_{p+1} = 2 - \frac{1}{a_p} = 2 - \frac{1}{\frac{p}{p-2}} = \frac{p}{p-1},
\]
which completes the induction and the proof of the proposition. \( \square \)

Without explicitly stating it, the following proposition was also given by Godrêche and Luck in [2]. We present it here, filling in the details.

**Proposition 6** ([2]). For \( n \geq 3 \) we have
\[
|A_n| = (n-1) \prod_{i=2}^{n-1} (n-i)^{f_{i-2}}.
\]

**Proof.** We give a proof by induction on \( n \). For the basis steps, \( n = 3 \) and \( n = 4 \) we have \( |A_3| = 2 \cdot 1^0 = 2 \) and \( |A_4| = 3 \cdot 2^0 \cdot 1^1 = 3 \). Now assume for induction that (6) holds for \( 4 \leq n \leq p \). Then from (5) and the induction assumption we have
\[
|A_{p+1}| = \frac{p}{p-1} |A_p||A_{p-1}|
\]
\[
= \frac{p}{p-1} \left( (p-1) \prod_{i=2}^{p-1} (p-i)^{f_{i-2}} \right) \left( (p-2) \prod_{i=2}^{(p-1)-1} ((p-1)-i)^{f_{i-2}} \right)
\]
\[
= p(p-2) \left( \prod_{i=2}^{p-1} (p-i)^{f_{i-2}} \right) \left( \prod_{i=2}^{(p-1)-1} (p-(i+1))^{f_{(i+1)-3}} \right)
\]
\[
= p(p-2) \left( \prod_{i=2}^{p-1} (p-i)^{f_{i-2}} \right) \left( \prod_{i=3}^{p-1} (p-i)^{f_{i-3}} \right)
\]
\[
= p(p-2)^{f_1}(p-2)^{f_0} \left( \prod_{i=3}^{p-1} (p-i)^{f_{i-2}+f_{i-3}} \right)
\]
\[
= p(p-1)^{f_0}(p-2)^{f_1} \left( \prod_{i=3}^{p-1} (p-i)^{f_{i-1}} \right)
\]
\[= p(p - 1)^f_0(p - 2)^f_1 \left( \prod_{i=4}^{p} (p - (i - 1))^{f_{(i-1)} - 1} \right)\]

\[= p \left( \prod_{i=2}^{(p+1)-1} ((p + 1) - i)^{f_{i-2}} \right),\]

which completes the induction. \(\blacksquare\)

The sequence \(\{|A_n|\}_{n=1}^{\infty}\) (see Table 1 on page 11) is the sequence A072042 in the On-line Encyclopedia of Integer Sequences [3].

From Proposition 6 we have

\[
\frac{\log |A_n|}{f_n} = \frac{1}{f_n} \log \left( n - 1 \prod_{i=2}^{n-1} (n - i)^{f_{i-2}} \right)
\]

\[= \frac{1}{f_n} \log(n - 1) + \sum_{i=2}^{n-1} \frac{f_{i-2}}{f_n} \log(n - i). \tag{7}\]

For large \(n\) the Fibonacci number \(f_n\) can be approximated by \(\frac{1}{\sqrt{5}} \left( \frac{1 + \sqrt{5}}{2} \right)^n\). Hence the sum in (7) converges when letting \(n\) tend to infinity and therefore clearly implies the existence of the limit. A numerical calculation in Maple gives

\[
\lim_{n \to \infty} \frac{\log |A_n|}{f_n} \approx 0.444399 \approx \log 1.559553.
\]

3 Prefixes

In this section we look at and present some properties of sets of prefixes of the inflated random Fibonacci words.

**Proposition 7.** For \(n \geq 3\) and \(k \geq 0\) we have

\[A_n[1, f_n - 1] = A_{n+k}[1, f_n - 1] \tag{8}\]

and symmetrically

\[A_n[2, f_n] = A_{n+k}[f_n - 2, f_n + 2] \tag{9}\]

**Proof.** We put our initial attention to equality (8). Let us first consider the case with \(k = 1\). We will give a proof by induction on \(n\). For the basis we have

\[A_3[1, 1] = \{0, 1\} = A_4[1, 1] \quad \text{and} \quad A_4[1, 2] = \{01, 10, 11\} = A_5[1, 2].\]
Now assume that (8) holds for $4 \leq n \leq m$. In the induction step $n = m + 1$ we get

$$A_{m+1}[1, f_{m+1} - 1] \subseteq (A_{m+1}A_m \cup A_mA_{m+1})[1, f_{m+1} - 1] = A_{m+2}[1, f_{m+1} - 1].$$

For the reversed inclusion let $x \in A_{m+2}[1, f_{m+1} - 1]$. By the recursion identity (2) we have to look at two cases. First, if $x \in (A_{m+1}A_m)[1, f_{m+1} - 1]$ then we clearly have $x \in A_{m+1}[1, f_{m+1} - 1]$ and we are done. Secondly, we have to deal with the case $x \in (A_mA_{m+1})[1, f_{m+1} - 1]$. For this case we use the induction assumption twice to obtain

$$A_{m-1}[1, f_{m-1} - 1] = A_m[1, f_{m-1} - 1] = (A_m[1, f_{m-1}])[1, f_{m-1} - 1] = (A_{m+1}[1, f_{m-1}])[1, f_{m-1} - 1] = A_{m+1}[1, f_{m-1} - 1].$$

By the help of the above equality it follows that

$$(A_mA_{m+1})[1, f_{m+1} - 1] = A_m(A_{m+1}[1, f_{m-1} - 1]) = A_m(A_{m-1}[1, f_{m-1} - 1]) = (A_mA_{m-1})[1, f_{m+1} - 1] \subseteq (A_mA_{m-1} \cup A_{m-1}A_m)[1, f_{m+1} - 1] = A_{m+1}[1, f_{m+1} - 1],$$

which completes the induction.

Now let us turn to the case with $k \geq 1$. We shall also prove this by induction. We proved the basis step, $k = 1$, above. Now we assume that (8) holds for $k = p$. For $k = p + 1$ we have by what we proved above

$$A_{n+p}[1, f_{n+p} - 1] = A_{n+p+1}[1, f_{n+p} - 1].$$

But then also $A_{n+p}[1, f_n - 1] = A_{n+p+1}[1, f_n - 1]$ and therefore by our induction assumption

$$A_n[1, f_n - 1] = A_{n+p}[1, f_n - 1] = A_{n+p+1}[1, f_n - 1],$$

which completes the proof of (8). Equality (9) is direct from (8) by symmetry.

**Proposition 8.** For $n \geq 4$ we have

$$A_n \subseteq (A_{n-1}[1, f_{n-1} - 1])\{0, 1\}^2(A_{n-2}[2, f_{n-2}]).$$

(10)
Proof. Denote the right hand side set in (10) by \( B_n \). If \( n = 4 \) we have \( B_4 = (A_3[1,1])\{0,1\}^2 = \{0,1\}^3 \), which clearly is a superset of \( A_4 \). If \( n \geq 5 \) then by Proposition 7 we have \( A_n[1,f_{n-1}] = B_n[1,f_{n-1}] \) and \( A_n[f_{n-1} + 2, f_n] = A_{n-2}[2, f_{n-2}] = B_n[f_{n-1} + 2, f_n] \). As \( B_n \) by definition consist of all words starting with a prefix in \( A_n[1,f_{n-1}] \) and ending with a suffix in \( A_n[f_{n-1} + 2, f_n] \) it must contain all sequences from \( A_n \).

Note that there clearly is a symmetric analogous of the result of Proposition 8.

4 Factor Sets

The aim of this section is to give a finite method for finding the factor set \( F_n \). For some small values of \( n \) we have

\[
F(A_3, f_2) = \{0,1\}, \quad F(A_4, f_3) = \{01,10,11\},
\]

\[
F(A_5, f_4) = \{001,010,011,100,101,110,111\},
\]

\[
F(A_6, f_5) = \left\{ \begin{array}{c}
00101,00110,00111,01001,01010, \\
01011,01100,01101,01110, \\
10010,10011,10100,10101,10110, \\
10111,11001,11010,11011,11100, \\
11101,11110
\end{array} \right\}.
\]

We shall shortly see that for \( n \geq 4 \) the above sets \( F(A_{n+1}, f_n) \) actually is \( F_n \). It is clear from above that \( F_4 \neq F(A_4, f_3) \) as words in \( F(A_5, f_4) \) have the sub-word 00, while this sub-word is not in \( F(A_4, f_3) \). By computer calculation we can find the size of \( F(A_{n+1}, f_n) \) for some additional small values of \( n \), see Table 1 on page 11.

Proposition 9. For \( k \geq 1 \) and \( n \geq 4 \) we have

\[
F(A_{n+1}, f_n) = F(A_{n+k}, f_n). \tag{11}
\]

Proof. We give a proof by induction on \( k \). For the basis step, \( k = 2 \), we see that \( F(A_{4+1}, f_4) = F(A_{4+2}, f_4) \), (this is direct as the word 000 is not a sub-word of an element in \( F(A_6, f_5) \)). Hence we may assume that \( n \geq 5 \). From the recursion (2) we have directly that \( F(A_{n+1}, f_n) \subset F(A_{n+2}, f_n) \). For the reversed inclusion we have by the recursion identity (2) and symmetry that it is enough to show that any sub-word of length \( f_n \) in \( A_{n+1}A_n \) is also a sub-word in \( A_{n+1} \). Let us define \( x_i = x[i,i+1] \) where \( x \in A_{n+1}A_n \) and \( 1 \leq i \leq f_{n+1} \).

For \( 1 \leq i \leq f_{n-1} + 1 \) we have directly that \( x_i \) is a sub-word in \( A_{n+1} \).

For \( f_{n-1} + 2 \leq i < 2f_{n-1} \) we have by the definition of \( x_i \) that it is a sub-word in \( A_{n+1}A_n[f_{n-1} + 2, f_{n+1} + f_{n-1} - 1] \). (The inequalities defining
the case are valid if \( f_{n-1} \geq 3 \), which requires that \( n \geq 5 \). By Proposition 7 and the recursion (2) we have

\[
(A_{n+1}A_n)[f_{n-1} + 2, f_{n+1} + f_{n-1} - 1] = \\
= (A_{n+1}[f_{n-1} + 2, f_{n+1}]) (A_n[1, f_{n-1} - 1]) \\
= (A_n[2, f_n]) (A_{n-1}[1, f_{n-1} - 1]) \\
= (A_nA_{n-1})[2, f_{n+1} - 1] \\
\subset (A_nA_{n-1} \cup A_{n-1}A_n)[2, f_{n+1} - 1] \\
= A_{n+1}[2, f_{n+1} - 1],
\]

which gives that \( x_i \) is a sub-word of \( A_{n+1} \).

For \( 2f_{n-1} \leq i \leq f_{n+1} \) we have that \( x_i \) is a sub-word in \( (A_{n+1}A_n)[f_n + 2, f_{n+2}] \). Proposition 7 and the recursion (2) now gives

\[
(A_{n+1}A_n)[f_n + 2, f_{n+2}] = (A_{n+1}[f_n + 2, f_{n+1}]) A_n \\
= (A_n[2, f_{n-1}]) A_n \\
= (A_nA_{n-1})[2, f_{n+1} - 1] \\
\subset (A_nA_{n-1} \cup A_{n-1}A_n)[2, f_{n+1} - 1] \\
= A_{n+1}[2, f_{n+1} - 1],
\]

and again we see that \( x_i \) is a sub-word of \( A_{n+1} \).

For finally \( i = f_{n+1} + 1 \) we have that \( x_i \) is an element in \( A_n \), and therefore also a sub-word in \( A_{n+1} \), which concludes the proof of the basis step.

Assume for induction that equality (11) holds for \( k = p \). For the induction step, \( k = p + 1 \), we have from what we just proved

\[
F(A_{n+p}, f_{n+p}) = F(A_{n+p+1}, f_{n+p}).
\]

But then we must also have \( F(A_{n+p}, f_n) = F(A_{n+p+1}, f_n) \). The induction assumption now gives \( F(A_n, f_n) = F(A_{n+p}, f_n) = F(A_{n+p+1}, f_n) \), which proves the proposition.

\[ \Box \]

**Proposition 10.** For \( n \geq 4 \) we have \( F_n = F(A_{n+1}, f_n) \).

**Proof.** It is clear that \( F(A_{n+1}, f_n) \subset F_n \). For the converse let \( x \in F_n \). Then there are two words \( u \) and \( v \) such that \( uxv \in A_m \) for some \( m \geq n + 1 \), that is \( x \in F(A_m, f_n) \). Proposition 9 now gives \( F(A_m, f_n) = F(A_{n+1}, f_n) \). \[ \Box \]

5 \hspace{1em} \textbf{Upper Bound}

The aim of this section is to give an upper bound of the size of \( |F_n| \) in terms of \( |A_n| \).

**Proposition 11.** For \( n \geq 3 \) and \( 1 \leq k \leq f_n - 1 \) we have

\[
|A_n[1, k]| \cdot |A_n[k + 1, f_n]| \leq 4^{n-2}|A_n|. \tag{12}
\]
Proof. It is clear from the symmetric structure of $A_n$ that we only have to give a proof for the case $1 \leq k \leq \lfloor \frac{1}{2} f_n \rfloor$. We give a proof by induction on $n$. For the basis case $n = 3$ we have

$$|A_3[1, 1]| \cdot |A_3[2, 2]| = |\{0, 1\}| \cdot |\{0, 1\}| = 2 \cdot 2 \leq 4 \cdot 2 = 4^{3-2}|A_3|.$$ 

Similarly we have for $n = 4$

$$|A_4[1, 1]| \cdot |A_4[2, 3]| \leq |A_4| \cdot |A_4| = 3 \cdot |A_4| \leq 4^{4-2}|A_4|$$

and for $n = 5$ with $k = 1, 2$

$$|A_5[k, 1]| \cdot |A_5[k + 1, 5]| \leq |A_5| \cdot |A_5| = 8 \cdot |A_5| \leq 4^{5-2}|A_5|.$$ 

Now assume for induction that (12) holds for $5 \leq n \leq p$. For the induction step, $n = p + 1$, we have from Proposition 7, Proposition 8, the induction assumption and 5 that

$$|A_{p+1}[1, k]| \cdot |A_{p+1}[k + 1, f_{p+1}]| = |A_{p}[1, k]| \cdot |A_{p+1}[k + 1, f_{p+1}]|$$

$$\leq |A_{p}[1, k]| \cdot \left( \left( A_{p}[1, f_{p} - 1] \right) \cdot \left( A_{p}[f_{p} - 1] \right) \right) |k + 1, f_{p+1}|$$

$$\leq |A_{p}[1, k]| \cdot |A_{p}[k + 1, f_{p}]| \cdot 2^2 \cdot |A_{p-1}[2, f_{p-1}]|$$

$$\leq 4^{p-2} \cdot |A_{p}| \cdot 4 \cdot |A_{p-1}|$$

$$\leq 4^{(p+1)-2}|A_{p+1}|,$$

which completes the induction.

The estimate in Proposition 11 seems to be far off from the true upper bound of the left hand side in (12). If we let

$$c_n = \max_{1 \leq k \leq f_n - 1} \frac{|A_n[1, k]| \cdot |A_n[k + 1, f_n]|}{|A_n|},$$

then computer calculations (see Table 1 on page 11) motivate the following conjecture

**Conjecture 12.** There is a constant $C$ such that for all $n \geq 3$ and $1 \leq k \leq f_n - 1$ we have

$$|A_n[1, k]| \cdot |A_n[k + 1, f_n]| \leq C \cdot |A_n|.$$ 

From our computer calculations (see Table 2 on page 11) we further conjecture that we may find $C < 3$, but we have to leave this question open as well as the more delicate question of finding the optimal value of such a constant, $c_n = \sup_{n \geq 3} c_n$.

We can now give the estimate we set out to find. Compared to computer calculations our estimate seems to be far from the optimal one, but it will be sufficient for our purpose.
Proposition 13. For \( n \geq 3 \) we have

\[
|F_n| \leq 2 \left( 4^{n-2} f_{n-1} + 1 \right) |A_n|.
\]

Proof. For \( n = 3 \) we have

\[
|F_3| = 4 \leq 2(4 \cdot 1 + 1) \cdot 2 = 2(4^{3-2} f_3^{-1} + 1)|A_3|.
\]

For \( n \geq 4 \) we have from Proposition 10 that \( F_n = F(A_n+1, f_n) \). Therefore it is enough to give an estimate of the number of subwords of length \( f_n \) that are in the words in \( A_{n+1} \). Let us start by considering the number of subwords in the subset \( A_n A_{n-1} \). We define \( x_k = x[k, k - 1 + f_n] \), where \( x \in A_n A_{n-1} \) for \( 1 \leq k \leq f_{n-1} + 1 \). Clearly we have \( x_k \in F_n \). If \( k = 1 \) then we have \( x_1 \in A_n \) and we see that there are \( |A_n| \) different \( x_1 \)'s. If \( 2 \leq k \leq f_{n-1} + 1 \) we have by Proposition 4

\[
x_k \in (A_n[k, f_n])(A_{n-1}[1, k - 1]) = (A_n[k, f_n])(A_n[1, k - 1]).
\]

Proposition 11 now gives that there are at most \( 4^{n-2}|A_n| \) different \( x_k \)'s for each of these \( k \)'s. Adding up the estimates and include the symmetric case, \( x \in A_{n-1} A_n \), completes the proof.

6 Proof of the Theorem

Finally, we have now gathered enough background to prove our theorem.

Proof of Theorem 3. We have already dealt with the existence of both limits in (3). Since clearly \( A_n \subset F_n \) it now follows from Proposition 13

\[
\lim_{n \to \infty} \frac{\log |A_n|}{f_n} \leq \lim_{n \to \infty} \frac{\log |F_n|}{f_n} \leq \lim_{n \to \infty} \frac{\log (2(4^{n-2} f_{n-1} + 1)|A_n|)}{f_n} \leq \lim_{n \to \infty} \frac{\log (4^{n-1} 2^n |A_n|)}{f_n} \leq \lim_{n \to \infty} \frac{3n \log 2}{f_n} + \lim_{n \to \infty} \frac{\log |A_n|}{f_n} = \lim_{n \to \infty} \frac{\log |A_n|}{f_n},
\]

which completes the proof.
7 Numerics

We present here the output from our computer calculations, which were obtained by a small JAVA-program on a standard PC.

| n  | fn  | An  | Fn  | F(A_{n+1}, fn) | cn     |
|----|-----|-----|-----|---------------|--------|
| 0  | 0   | 0   | 0   |               |        |
| 1  | 1   | 1   | 2   | 1             |        |
| 2  | 1   | 1   | 2   | 2             |        |
| 3  | 2   | 2   | 4   | 3             | 2.0    |
| 4  | 3   | 3   | 7   | 7             | 2.0    |
| 5  | 5   | 8   | 22  | 22            | 2.0    |
| 6  | 8   | 30  | 108 | 108           | 2.13333|
| 7  | 13  | 288 | 1356| 1356          | 2.11111|
| 8  | 21  | 10080| 65800| 65800     | 2.17143|
| 9  | 34  | 3317760| 30139200| 30139200| 2.16389|

8 Acknowledgment

The author wishes to thank Michael Baake, Franz Gähler and Claudia Lütkehöller at Bielefeld University, Germany and Tomas Persson at the Polish Academy of Sciences, Warsaw, Poland for presenting the problem, our discussions of it and for reading drafts of the manuscript. This work was partly supported by the German Research Council (DFG), via CRC 701.

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