Analog quantum error correction

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Abstract: Quantum error-correction routines are developed for continuous quantum variables such as position and momentum. The result of such analog quantum error correction is the construction of composite continuous quantum variables that are largely immune to the effects of noise and decoherence.

The quantum systems used for quantum computation and quantum communications are small, sensitive, and easily perturbed [1-8]. The theory of quantum error-correcting codes provides a new set of techniques for protecting quantum systems against the effects of noise and decoherence [9-29]. Conventional quantum error-correcting codes are only effective for discrete variables, however. This letter presents a set of analog quantum error-correcting routines that protect continuous variables such as position and momentum against noise and decoherence. These error-correcting routines can in principle be enacted using simple Hamiltonian operations to stabilize the states of arbitrary continuous quantum variables. Particular applications include error-correction for quantum communications using continuous
variables such as photon momentum, and for analog quantum computers used for simulating continuous quantum systems [30-31].

First consider the problem of correcting errors in a classical discrete system. The simplest binary error-correcting routine is triple modular redundancy, in which three bits are initially set to the same value and checked at regular intervals to see if they still have the same value: if one of them differs, it is reset to the value of the two others. If the error rate per bit per unit time is $\lambda$, then performing this ‘voting’ routine at intervals of time $\delta t << 1/\lambda$ results in a new error rate of $3\lambda^2 \delta t << \lambda$.

The discrete error-correcting technique of triple modular redundancy can be adapted simply to continuous classical variables. Consider three continuous variables $x_1 x_2 x_3$, initially set to the same value $x$. If at some brief time later one of the three is found to differ from the other two, it is reset to the majority value, $M$. (If all three differ, then $M$ can be taken to be the average value of the two variables that differ the least.) The resetting can be accomplished by a simple nonlinear dynamics such as $\ddot{x}_j = -k(x_j - M) - \gamma \dot{x}_j$, where $\gamma, \sqrt{k} >> \lambda$. Just as in the discrete case, if the probability of a variable being perturbed per unit time is $\lambda$, then performing this nonlinear ‘continuous voting’ routine at intervals of time $\delta t << 1/\lambda$ results in a new error rate of $3\lambda^2 \delta t << \lambda$, which can be made arbitrarily small by reducing $\delta t$.

This classical continuous error-correcting routine is clearly dissipative and decohering. It can be modified to preserve quantum coherence, however, and can be used in a quantum context to protect against some forms of quantum error. Consider three continuous ‘position’ quantum variables with states $|x_1 x_2 x_3\rangle_{123}$, and errors corresponding to unitary operators $e^{-iQ(P_j)}$, where $P_j = -i\partial/\partial x_j$ is the ‘momentum’ operator on the $j$’th variable and $Q$ is a polynomial function of $P_j$ (we call these variables position and momentum for convenience only: the method works for any continuous variable and its conjugate). Such an error takes

$$|x\rangle_i \rightarrow e^{-iQ(P_j)}|x\rangle_j = (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{-ipx - iQ(p)}|p\rangle_j dp$$  \hspace{1cm} (1)

where $|p\rangle_j = (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{ipx}|x\rangle_j dx$. The error acts on only one variable: $|x\rangle_k \rightarrow |x\rangle_k$ for $k \neq j$. For example, $Q(P_j) = \delta x P_j$ takes

$$|x\rangle_j \rightarrow (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{-ipx - ip\delta x}|p\rangle_j dp = |x + \delta x\rangle_j.$$ \hspace{1cm} (2)

To correct for these errors, apply the following quantum ‘continuous voting’ procedure. We assume that a variable can be prepared in the state $|0\rangle_j$ by some
dissipative process such as cooling, and that the state $|x\rangle_j$ can also be prepared, e.g., by applying the ‘displacement’ Hamiltonian $\eta x P_j$ to the state $|0\rangle_j$ for a time $1/\eta$.

To ‘vote,’ apply the following procedure to three continuous quantum variables, initially in the state $|xxx\rangle_{123}$, together with three ancilla variables $|x_1x_2x_3\rangle_{1'2'3'}$, initially in the state $|000\rangle_{1'2'3'}$:

0) Suppose that an error occurs to one of the variables, e.g., the second one:

$$|x\rangle_2 \rightarrow e^{-iQ(P_2)}|x\rangle_2 = (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{-ipx-iQ(p)}|p\rangle_2 dp = \int_{-\infty}^{\infty} \alpha(x,x')|x\rangle_2 dx'$$

(3) where $\alpha(x,x') = (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{-ip(x-x')-iQ(p)}dp$. Reprepare the ancilla variables in the state $|000\rangle_{1'2'3'}$ (this corrects any error that has occurred to the ancillae). The overall state of the variables and the ancillae is now

$$\left(|x\rangle_1|0\rangle_1\right) \left(\int_{-\infty}^{\infty} \alpha(x,x')|x\rangle_2 dx' |0\rangle_2\right) \left(|x\rangle_3|0\rangle_3\right)$$

(4)

1) Perform a continuous quantum analog of voting. We will assume that we can perform simple real-number operations such as comparing the values of two variables to see if they are equal, and adding the value of one variable to another. So for example, we will assume that we can perform operations such as comparing $|x_1\rangle_1$ and $|x_2\rangle_2$ to see if $x_1 = x_2$, and if they are, performing operations such as $|x_1\rangle_1|x_2\rangle_2|x_3\rangle_3 \rightarrow |x_1\rangle_1|x_2\rangle_2|x_3 + x_1\rangle_3$. Such operations are reversible and correspond to unitary transformations on Hilbert space. They can be accomplished by the application of simple interactions between variables. For example, the conditional addition operation just described can be accomplished by applying the Hamiltonian $\eta \Delta(X_1, X_2) X_1 P_3$ for time $1/\eta$, where

$$\Delta(X_1, X_2) = \int_{-\infty}^{\infty} |x\rangle_1\langle x| \otimes |x\rangle_2\langle x| dx.$$
equal. So one by one, compare each of the $|x\rangle_i$ to the other two. Let $ijk$ be some permutation of 123. If $x_i = x_j = x$, then add the value of $x_k - x$ to the ancilla state $|0\rangle_k'$. If $x_i \neq x_j$ then do nothing. In our case, only $|x\rangle_2$ and its ancilla will be affected. The variables and ancillae are now in the state

$$\langle x\rangle_1|0\rangle_1' \left( \int_{-\infty}^{\infty} \alpha(x, x') |x\rangle_2|x'\rangle_2| x - x'\rangle_2^d x' dx' \right) (|x\rangle_3|0\rangle_3'). \tag{5}$$

(2) Now if $x_i = x_j$ subtract the value of the $k$’th ancilla variable from the original $k$’th variable, leaving the state

$$\langle x\rangle_1|0\rangle_1' \left( \int_{-\infty}^{\infty} \alpha(x, x') |x - x'\rangle_2^d x' dx' \right) (|x\rangle_3|0\rangle_3'). \tag{6}$$

Substituting in the explicit expression for $\alpha(x, x')$ given above allows this state to be written as

$$\langle x\rangle_1|0\rangle_1' \left( \int_{-\infty}^{\infty} e^{-ip(x-x') - iQ(p)} |x' - x\rangle_2^d x' dp \right) (|x\rangle_3|0\rangle_3').$$

$$= \langle x\rangle_1|0\rangle_1' \left( |x\rangle_2 (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{-iQ(p)} |p\rangle_2 dp \right) (|x\rangle_3|0\rangle_3'). \tag{7}$$

$$= \langle x\rangle_1|0\rangle_1' \left( |x\rangle_2 e^{-iQ(P_2')} |0\rangle_2' \right) (|x\rangle_3|0\rangle_3').$$

The error has now been corrected.

This procedure corrects the error by restoring the three variables to the original continuous ‘codeword’ $|xxx\rangle_{123}$ while leaving the ancilla in a state that is independent of the initial value of $x$. The fact that the ancilla is in a state that depends only on error operator $e^{-iQ(P_i)}$ applied and not on the particular ‘codeword’ to which it is applied means that the procedure restores not only continuous codewords but arbitrary superpositions of the codewords $\int_{-\infty}^{\infty} \psi(x)|xxx\rangle dx$.

To continue correcting errors, simply return the ancilla variables to $|000\rangle_{1'2'3'}$ and apply the procedure again a time $\delta t$ later. Just as in the classical case, performing the error-correcting routine at intervals $\delta t$ reduces the error rate from $\lambda$ to $3\lambda^2 \delta t$, which can be made as small as desired by decreasing $\delta t$.

It can easily be seen by interchanging the roles of $x$ and $p$ above that continuous codewords of the form $|ppp\rangle_{123}$ can be protected against arbitrary errors of the form $e^{iR(X_j)}$ where $X_j$ is the position operator on the $j$’th variable and $R$ is a polynomial function of $X_j$. In analogy to the $|xxx\rangle$ error-correcting routine, we assume that variables and ancillae can be prepared in momentum eigenstates, $|p = 0\rangle_j$, and that states $|p\rangle_j$ can be created by applying the ‘boost’ Hamiltonian $\eta p X_j$ to the state
\( |p = 0\rangle_j \) for a time \( 1/\eta \). The ancillae for the \( |ppp\rangle \) error-correcting routine are be prepared initially in the state \( |p = 0\rangle_{j'} \), or they can be prepared in position eigenstates \( |x = 0\rangle_{j'} \) as before, and the interactions between variables and ancillae adjusted to convert the value of momentum in a variable to the value of position in the ancilla.

The following algorithm corrects both phase and displacement errors. Define the state

\[
|p\rangle_{123} \equiv (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{i px} |xxx\rangle_{123} dx.
\]  

(Such a state can be created from the state

\[
|p\rangle_{1} |0\rangle_{2} |0\rangle_{3} \equiv (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{i px} |x\rangle_{1} |0\rangle_{2} |0\rangle_{3} dx
\]

by applying the Hamiltonian \( \eta X_j P_j \) for time \( 1/\eta \) to effect the unitary operation

\[
|x\rangle_{1} |y\rangle_{j} \rightarrow |x\rangle_{1} |x + y\rangle_{j}
\]

for \( j = 2, 3 \).

The error operator \( e^{i R(X_j)} \) has the same effect on the triple-variable state \( |p\rangle_{123} \) that it has on the single-variable state \( |p\rangle_j \):

\[
|p\rangle_{123} \rightarrow e^{i R(X_j)} |p\rangle_{123} = (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{i px + iR(x)} |xxx\rangle_{123} dx.
\]  

This error can be corrected in an analogous way to the errors on single continuous variables: create redundant states of the nine variables \( |p_{123}p_{456}p_{789}\rangle_{1...9} \) together with a set of three ancilla variables originally in the state \( |000\rangle_{ABC} \), where ancilla variable \( A \) is used as the ancilla for the triple of variables 123, \( B \) is used for 456, and \( C \) is used for 789, then carry out the same error-correcting dynamics as above, but as a function of the continuous variables \( p \) that label the states \( |p\rangle \). That is, phase errors on the triply-redundant state of triply-redundant continuous variables can be corrected by applying essentially the same error-correcting routine as before.

To correct any combination of phase and displacement errors on one variable, first apply the \( |xxx\rangle \) error-correction routine for error operators of the form \( e^{-iQ(P_j)} \) to each of the three triples of variables, \( 123, 456, 789 \), then apply the \( |ppp\rangle \) error-correction routine for error operators of the form \( e^{i R(X_j)} \) to the nine variables as a whole. The basic idea of this continuous quantum error-correcting routine is the same as Shor’s binary quantum error correcting routine [9]: using triple modular redundancy twice (‘triple-triple’ modular redundancy) corrects both phase and displacement errors. This sequence of error correcting steps compensates for the effect
of any error operator of the form $e^{-iQ(X_j,P_j)}$, where $Q(X_j,P_j)$ is now a polynomial in the operators $X_j, P_j$.

To see the error-correction explicitly, use the commutation relation $[X_j, P_j] = i$ to write $e^{-iQ(X_j,P_j)} = \sum_{m,n \geq 0} q_{mn} P_j^n X_j^m$. Look at what happens when an error of this form occurs to one of the variables, for example, the first ($j=1$). We have

$$|p_{123}p_{456}p_{789}\rangle_{1...9}|0\ldots0\rangle_{1'...9'}|000\rangle_{ABC}$$

$$\rightarrow \sum_{mn} q_{mn} P_1^m X_1^n (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{ipx} |xx\rangle_{123} dx$$

(10)

$$|p_{456}|p_{789}|0\ldots0\rangle_{1'...9'}|000\rangle_{ABC}$$

which can be rewritten using the decompositions $|x\rangle = (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{-ipx} |p\rangle dp$, $|p\rangle = (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} e^{-ipx'} |x'\rangle dp'$, as

$$\sum_{mn} q_{mn}(1/\sqrt{2\pi})^3 \int_{-\infty}^{\infty} p^m x^n e^{ipx} e^{-ip(x-x')} |x'\rangle_1 |xx\rangle_{23} dx dx' dp$$

(11)

$$|p_{456}|p_{789}|0\ldots0\rangle_{1'...9'}|000\rangle_{ABC}.$$  

Now proceed as before, comparing $x_1, x_j, x_k$, and if $x_i = x_j = x$, adding $y = x' - x$ to the value of the ancilla state $|0\rangle_k$ and subtracting $x' - x$ from the value of the state $|x'\rangle_k$. Only the first variable and its ancilla state will be affected, resulting in the state

$$\sum_{mn} q_{mn}(1/\sqrt{2\pi})^3 \int_{-\infty}^{\infty} p^m x^n e^{ipx} e^{-ipy} |y\rangle_1 |xx\rangle_{123} dx dy dp$$

$$|p_{456}|p_{789}|0\ldots0\rangle_{1'...9'}|000\rangle_{ABC}$$

(12)

$$\sum_{m,n} q_{mn}(1/\sqrt{2\pi}) \int_{-\infty}^{\infty} x^n e^{ipx} |xx\rangle_{123} dx$$

$$|p_{456}|p_{789} P_1^m |0\rangle_{1'} |0\ldots0\rangle_{2'...9'}|000\rangle_{ABC},$$

where $P_1^m$ acts only on the first ancilla variable. The error-correction routine for states of the form $|xx\rangle$ has transferred the effect of the $P_j^m$ part of the the error operator from the codeword to the ancilla.

Similarly, applying the $|ppp\rangle$ error-correction to the state in (12) transfers the effect of the $X_j^n$ part of the error-operator from the codeword to the ancilla, resulting in the state

$$e^{ipx} |xx\rangle_{123} dx |p_{456}|p_{789}$$

$$\left( \sum_{m,n} q_{mn} P_1^m |0\rangle_{1'} X_A^n |0\rangle_{A} |0\ldots0\rangle_{2'...9'}|00\rangle_{BC} \right) \rightarrow |p_{123}|p_{456}|p_{789}|e^{-iQ(X_A,P_1')}|0\ldots0\rangle_{1'...9'}|000\rangle_{ABC}.$$  

(13)
The error has now been corrected. The ancillae can be reset and the procedure repeated to provide ongoing error correction.

To summarize: In each term of the polynomial expansion of the error operator, the application of $|xxx\rangle$ error-correcting routine to the triple of continuous variables containing $j$ restores the triple where the error occurred to a superposition of the form $\int_{-\infty}^{\infty} \beta_n(p,p')|p'\rangle dp'$, where $\beta_n(p,p') = (1/\sqrt{2\pi}) \int_{-\infty}^{\infty} x^n e^{i(p-p')x} dx$. The subsequent application of the $|ppp\rangle$ error-correction routine to the triple of triples then restores the nine variables as a whole to the state $|ppp\rangle_{1...9}$. The fact that the state of the ancillae after each error-correcting routine depends only on what errors occurred and not on which codeword $|ppp\rangle_{1...9}$ the system was in implies that arbitrary superpositions of the form $\int_{-\infty}^{\infty} \psi(p)|ppp\rangle_{1...9} dp$ are also restored by the continuous error-correction routine.

The analog quantum error-correcting routine presented above corrects for errors that are arbitrary polynomials in $X_j$ and $P_j$. It can be enacted in principle using simple operations on the real numbers such as comparing and adding two numbers. What happens when these operations can only be performed to finite precision? By going through the error-correcting routine and following what happens when comparison and addition are performed to finite precision $\delta$, one can verify that the procedure still works as long as (i) the wave-function $\psi$ does not vary significantly over scales $\delta$, and (ii) the expectation values for the error operators on the range of $\psi$ do not vary significantly over scales $\delta$. Perhaps the easiest way to see why such inexact error-correction still works is to note that when (i-ii) hold for finite precision $\delta$ in manipulations of continuous variables the system behaves like an infinite-dimensional discrete system with states $|x_n\rangle = |n\delta\rangle$. The continuous error-correcting scheme above, performed at finite precision, still functions as an error-correcting scheme for the discrete infinite-dimensional system. Similarly, the method described here generalizes in a straightforward fashion to systems that are continuous in one variable and discrete in the complementary variable (e.g., a particle in a box).

Our method is continuous but time-dependent: it may be possible to devise error-correcting dynamics that are time-independent as well. A particularly interesting application of analog quantum error correction is to a collection of systems that evolve according to a master equation:

$$\dot{\rho} = -i[H,\rho] + \sum_m (L_m\rho L_m^\dagger - (1/2)L_m^\dagger L_m\rho - (1/2)\rho L_m^\dagger L_m)$$

Suppose that each $L_m$ is a polynomial function of position and momentum opera-
tors acting on individual subsystems (i.e., each subsystem sees an effectively distinct environment; when the particles are coupled to the same environment it may be possible to use symmetrized states of the particles to resist noise and decoherence [32]). This is the typical case of particles each governed by a distinct single particle master equation such as the optical master equation. Inducing the proper interactions with ancillary continuous systems and applying the error-correcting routine given above at intervals $\delta t$ allows nine subsystems to be grouped together into one composite system whose states $\int_{-\infty}^{\infty} \psi(p)|ppp\rangle_{1\ldots9}dp$ are unaffected by the dynamics (14) to first order in $\delta t$. For the continuous error-correction to be effective, it must be repeated at intervals shorter than the dynamic time-scales of the system such as its decoherence time or spontaneous emission time. The analog quantum error-correcting routine presented here allows the creation of joint states of a composite continuous system that are largely immune to the effects of interference and noise in principle. In practice, of course, performing the continuous ‘quantum logic gates’ necessary to enact the analog error-correcting scheme is likely to prove difficult.

We have presented a quantum error-correcting routine for continuous variables. The routine allows the creation of states of a composite system that resist the effects of errors and noise. For simplicity of exposition, we presented a method for analog quantum error correction based on Shor’s original error-correcting routine for qubits. A variety of other continuous quantum error-correcting routines can be constructed based on other discrete quantum codes. In particular, in analogy to [29], it should be possible to devise a ‘perfect’ analog quantum error-correcting code using only five continuous variables, although the dynamics of the error correction will be more complicated than the simple continuous voting used here [33]. The quantum error-correcting mechanism described here is an example of a feedback loop that preserves quantum coherence as proposed by Lloyd [34]. The nonlinear dynamics cause the ancilla variables to become correlated with the system in a coherent manner, and the information that they possess is used coherently to restore the system to its desired state.

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