ALGORITHMIC CONSTRUCTIONS OF UNITARY MATRICES AND TIGHT FRAMES

JANET C. TREMAIN

Abstract. We give a number of algorithms for constructing unitary matrices and tight frames with specialized properties. These were produced at the request of researchers at the Frame Research Center (www.framerc.org) to help with their research on fusion frames, the Kadison-Singer Problem and equiangular tight frames.

1. Introduction

We will give a sequence of algorithms for constructing unitary matrices (i.e. orthonormal bases) and tight frames with specialized properties. For unitary matrices, we just need to construct square matrices with rows and columns orthogonal and row and column numbers square summing to one. For tight frames, we need to produce \( M \times N \) matrices with \( M \geq N \) which have orthogonal columns and whose column numbers square sum to a fixed constant. Then the row vectors of this matrix will form a tight frame for \( \mathbb{H}_N \). We will be interested in a number of special properties such as sparse orthonormal bases for hyperplanes, matrices with first row a constant, first two rows a constant, matrices made from just 2 constants, matrices built out of unitary matrices, matrices with a large portion of the norms of their vectors allocated to special positions in the matrix etc. We will work specifically with 2-tight frames since it is known that paving these is equivalent to solving the famous Kadison-Singer Problem \([1, 4, 5]\).

Recall that a family of vectors \( \{f_i\}_{i=1}^M \) is a frame for a Hilbert space \( \mathbb{H}_N \) if there are constants \( 0 < A \leq B < \infty \), called the lower and upper frame bounds respectively so that

\[
A\|f\|^2 \leq \sum_{i=1}^M |\langle f, f_i \rangle|^2 \leq B\|f\|^2.
\]

The frame is \( A \)-tight if \( A = B \) and Parseval if \( A = B = 1 \). It is equal norm of \( \|f_i\| = c \), for all \( i = 1, 2, \ldots, M \), and unit norm if \( c = 1 \). The analysis operator of the frame is \( T : \mathbb{H}_N \to \ell_2^M \) given by

\[
Tf = \sum_{i=1}^M \langle f, f_i \rangle e_i,
\]

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where \( \{e_i\}_{i=1}^M \) is the natural orthonormal basis of \( \ell_2^M \). The adjoint of the analysis operator is the \textit{synthesis operator} which satisfies

\[
T^* \left( \{a_i\}_{i=1}^M \right) = \sum_{i=1}^M a_i f_i.
\]

The frame operator is the positive, self-adjoint invertible operator given by

\[
S =: T^* T f = \sum_{i=1}^M \langle f, f_i \rangle f_i.
\]

The synthesis operator of the frame is the matrix with the frame vectors as column vectors:

\[
A = \begin{bmatrix}
| & | & | & | & | \\
| & | & | & | & | \\
| & | & | & | & | \\
\end{bmatrix}
\]

2. Sparse Orthonormal Bases for Hyperplanes

For their work on \textit{sparsity} the researchers at the Frame Research Center needed the sparsest orthonormal bases for hyperplanes in \( \mathbb{R}^N \) which do not contain the span of any subset of the orthonormal basis. We believe that this class of hyperplanes is given by the following cases - although at this time we do not have a proof that these are the sparsest.

\textbf{dimension n=2:}

\[
A_2 = \begin{bmatrix} 1 & 1 \end{bmatrix}
\]

\textbf{dimension n=3:}

\[
\begin{bmatrix} 1 & 2 & 1 \\ 0 & -1 & 2 \end{bmatrix}
\]

\textbf{dimension n=4:}

\[
A_4 = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 0 & 0 \\ 0 & 0 & 1 & -1 \end{bmatrix}
\]

The above has sparsity 8 and it is easily checked that any orthonormal basis for this hyperplane in \( \mathbb{R}^4 \) has sparsity \( \geq 8 \).

\textbf{dimension n=5:} Here we combine the earlier cases of \( A_2 \) and \( A_3 \) and then add a non-zero row on top.

\[
A_5 = \begin{bmatrix} * & * & * & * & * \\ * & * & 0 & 0 & 0 \\ 0 & * & 0 & 0 & * \\ 0 & 0 & 0 & * & * \end{bmatrix}
\]

The above has sparsity 11.
dimension \( n=6 \): Here, we will combine the cases \( 2n = 2 \) with \( 2n = 4 \) and then insert a new top row.

\[
A_6 = \begin{bmatrix}
1 & 1 & 1 & 1 & 1 & 1 \\
1 & 1 & -1 & -1 & 0 & 0 \\
1 & -1 & 0 & 0 & 0 & 0 \\
0 & 0 & 1 & -1 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & -1
\end{bmatrix}
\]

The above has sparsity 16.

dimension \( n=7 \): Here we combine the earlier cases of \( n = 4 \) and \( n = 3 \).

\[
A_7 = \begin{bmatrix}
* & * & * & * & * & * & * \\
* & * & * & 0 & 0 & 0 & 0 \\
* & * & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & * & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & * & * & * \\
0 & 0 & 0 & 0 & * & * & *
\end{bmatrix}
\]

This has sparsity 20.

dimension \( n=8 \): In this case we combine two of the cases \( 2n = 4 \) and add a new top row.

\[
A_8 = \begin{bmatrix}
1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\
1 & 1 & -1 & -1 & 0 & 0 & 0 & 0 \\
1 & -1 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 1 & -1 & 0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 1 & 1 & -1 & -1 \\
0 & 0 & 0 & 0 & 1 & -1 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 1 & -1
\end{bmatrix}
\]

This has sparsity 24.

The general cases here are:

dimension \( 2n \) with \( n \) even: We combine two of the cases \( A_{2n} \) and add a new top row.

\[
A_{4n} = \begin{bmatrix}
2n \ \text{ones} & 2n \ \text{ones} \\
A_{2n} & A_{2n}
\end{bmatrix}
\]

This has sparsity:

\[
2[\text{sparsity of } A_{2n}] + 4n.
\]

dimension \( 2(n+1) \), \( n+1 \) odd: We combine the previous cases \( n+2 \) and \( n \).

\[
\begin{bmatrix}
\begin{array}{cc}
n+2 & \text{ones} \\
A_{n+2} & A_n
\end{array}
\end{bmatrix}
\]
This has sparsity

\[ \text{sparsity of } A_{n+2} + \text{sparsity of } A_n + 2n + 2. \]

**Dimension** \( 2n + 1 = 2^k + 1 + 2i \) with \( i \) even: Here we combine the cases \( n = 2^{k-1} + 2i \) and \( 2^{k-1} + 1 \) and add a new top row.

\[
A_{2n+1} = \begin{bmatrix}
2^{k-1} + 2i \text{ stars} & 2^{k-1} + 1 \text{ stars} \\
A_{2k-1+2i} & A_{2k-1+1}
\end{bmatrix}
\]

**Dimension** \( 2n + 1 = 2^k + 1 + 2i \) with \( i \) odd:

\[
A_{2n+1} = \begin{bmatrix}
2^{k-1} \text{ stars} & 2^{k-1} + 1 + 2i \text{ stars} \\
A_{2k-1+2i} & A_{2k-1+1+2i}
\end{bmatrix}
\]

3. Unitary Matrices with First Row a Constant

The researchers at the Frame Research Center needed unitary matrices with the first row a constant and maximal sparsity for their work on equian-
gular fusion frames. They also wanted a large number of zero’s in the matrix. For this algorithm, we can just alter the top row, then add a new top row and first column to the previous example.

To make this more clear, we do two small examples and then the general case. First we do \( 3 \times 3 \).

\[
A_3 = \begin{bmatrix}
\frac{1}{\sqrt{3}} & \frac{1}{\sqrt{3}} & \frac{1}{\sqrt{3}} \\
\frac{-\sqrt{3}}{2} & \frac{1}{\sqrt{3}} & \frac{1}{\sqrt{2}} \\
0 & \frac{-1}{\sqrt{2}} & \frac{1}{\sqrt{2}}
\end{bmatrix}
\]

Next we do \( 4 \times 4 \).

\[
A_4 = \begin{bmatrix}
\frac{1}{\sqrt{3}} & \frac{1}{\sqrt{4}} & \frac{1}{\sqrt{4}} & \frac{1}{\sqrt{4}} \\
\frac{-\sqrt{3}}{\sqrt{4}} & \frac{1}{\sqrt{4}} & \frac{1}{\sqrt{3}} & \frac{1}{\sqrt{4}} \\
\frac{0}{\sqrt{3}} & \frac{-\sqrt{3}}{\sqrt{3}} & \frac{1}{\sqrt{3}} & \frac{1}{\sqrt{3}} \\
0 & 0 & \frac{-1}{\sqrt{2}} & \frac{1}{\sqrt{2}}
\end{bmatrix}
\]

Now for the general \( n \times n \) case.
4. Unitary Matrices with First Two Rows a Constant Modulus

The researchers at the FRC needed these matrices for their work on constructing tight fusion frames.

Again, we give a few small cases to get this started.

**dimension 4:**

\[
A_4 = \begin{bmatrix}
-1 & 1 & 1 & 1 \\
1 & -1 & 1 & 1 \\
1 & 1 & -1 & 1 \\
1 & 1 & 1 & -1 \\
\end{bmatrix}
\]

**dimension 6:**

\[
A_6 = \begin{bmatrix}
\sqrt{\frac{1}{8}} & \sqrt{\frac{1}{8}} & \sqrt{\frac{1}{8}} & \sqrt{\frac{1}{8}} & \sqrt{\frac{1}{8}} & \sqrt{\frac{1}{8}} \\
-\sqrt{\frac{1}{8}} & -\sqrt{\frac{1}{8}} & -\sqrt{\frac{1}{8}} & \sqrt{\frac{1}{8}} & \sqrt{\frac{1}{8}} & \sqrt{\frac{1}{8}} \\
-\sqrt{\frac{3}{8}} & \sqrt{\frac{1}{8}} & \sqrt{\frac{1}{8}} & 0 & 0 & 0 \\
0 & \sqrt{\frac{1}{2}} & \sqrt{\frac{1}{2}} & 0 & 0 & 0 \\
0 & 0 & 0 & -\sqrt{\frac{3}{8}} & \sqrt{\frac{1}{8}} & \sqrt{\frac{1}{8}} \\
0 & 0 & 0 & 0 & \sqrt{\frac{1}{2}} & \sqrt{\frac{1}{2}} \\
\end{bmatrix}
\]

For \(A_8\) we can use a Hadamard again. But I was asked for "sparse" matrices doing this.

**dimension 8:**
And in general:

\[ A_{2n} = \]

\[
\begin{pmatrix}
\sqrt{\frac{1}{2n}} & \sqrt{\frac{1}{2n}} & \cdots & \sqrt{\frac{1}{2n}} & \sqrt{\frac{1}{2n}} & \cdots & \sqrt{\frac{1}{2n}} \\
-\sqrt{\frac{1}{2n}} & -\sqrt{\frac{1}{2n}} & \cdots & -\sqrt{\frac{1}{2n}} & -\sqrt{\frac{1}{2n}} & \cdots & -\sqrt{\frac{1}{2n}} \\
-\sqrt{\frac{n-1}{n}} & -\sqrt{\frac{n-1}{n(n-1)}} & \cdots & -\sqrt{\frac{n-1}{n}} & -\sqrt{\frac{n-1}{n(n-1)}} & \cdots & -\sqrt{\frac{n-1}{n}} \\
\vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\
0 & 0 & \cdots & -\sqrt{\frac{1}{2}} & 0 & \cdots & 0 \\
0 & 0 & \cdots & 0 & -\sqrt{\frac{n-1}{n}} & \cdots & \sqrt{\frac{1}{n(n-1)}} \\
\vdots & \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\
0 & 0 & \cdots & 0 & 0 & \cdots & \sqrt{\frac{1}{2}}
\end{pmatrix}
\]

5. Unitary Matrices: Made from 2 Constants

We will look at various ways to construct unitary matrices from just two constants.

5.1. First Case. In this section we construct a very special class of unitary matrices of the form

\[
\frac{1}{n} \begin{pmatrix}
-b & a & a & a & \cdots & a \\
a & -b & a & a & \cdots & a \\
a & a & -b & a & \cdots & a \\
\vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\
a & a & a & a & \cdots & -b
\end{pmatrix}
\]
For the case of $\mathbb{R}^3$ this looks like:

$$\frac{1}{3} \begin{bmatrix} -1 & 2 & 2 \\ 2 & -1 & 2 \\ 2 & 2 & -1 \end{bmatrix}$$

For the case $\mathbb{R}^6$ this looks like:

$$\frac{1}{3} \begin{bmatrix} -2 & 1 & 1 & 1 & 1 & 1 \\ 1 & -2 & 1 & 1 & 1 & 1 \\ 1 & 1 & -2 & 1 & 1 & 1 \\ 1 & 1 & 1 & -2 & 1 & 1 \\ 1 & 1 & 1 & 1 & -2 & 1 \\ 1 & 1 & 1 & 1 & 1 & -2 \end{bmatrix}$$

Or in a general form, the above looks like:

$$\frac{1}{n} \begin{bmatrix} -2 & 2 & 2 & 2 & 2 & 2 \\ 2 & -2 & 2 & 2 & 2 & 2 \\ 2 & 2 & -2 & 2 & 2 & 2 \\ 2 & 2 & 2 & -2 & 2 & 2 \\ 2 & 2 & 2 & 2 & -2 & 2 \\ 2 & 2 & 2 & 2 & 2 & -2 \end{bmatrix}$$

And for the general case of $\mathbb{R}^n$ we have:

$$\frac{1}{n} \begin{bmatrix} -(n-2) & 2 & 2 & \cdots & 2 \\ 2 & -(n-2) & 2 & \cdots & 2 \\ 2 & 2 & -(n-2) & \cdots & 2 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 2 & 2 & 2 & \cdots & -(n-2) \end{bmatrix}$$

5.2. **Second Case.** In this section we construct a class of unitary matrices of the form

$$\begin{bmatrix} -a & a & a & -b & b & b \\ a & -a & a & b & -b & b \\ a & a & -a & b & b & -b \\ a & a & a & -a & b & b \\ -b & b & b & a & -a & -a \\ b & -b & b & -a & a & -a \\ b & b & -b & -a & a & -a \\ b & b & b & -b & -a & a \end{bmatrix}$$

The rows and columns of this matrix are clearly orthogonal. To make the rows and columns square sum to one, we just need to have that $4a^2 + 4b^2 = 1$. For example, in the above case, $b = \frac{1}{\sqrt{20}}$ and $a = 2b$ work. The above matrix represents a generalization of the Hadamard construction. In general it looks like:
Theorem 5.1. Let $A, B$ be $n \times n$ unitary matrices. Choose $a, b$ so that $a^2 + b^2 = 1$. Then the following matrix is unitary:

$$C = \begin{bmatrix} aA & bB \\ aA & -bB \end{bmatrix}$$

Proof. A direct calculation gives that $C^*C = I$. \qed

More general than this is:

Theorem 5.2. Let $A_1, A_2, \ldots, A_n$ be $m \times m$ unitary matrices and let $(a_{ij})_{i,j=1}^n$, be a unitary matrix.

Then the following matrix is a $nm \times nm$ unitary matrix:

$$B = \begin{bmatrix} a_{11}A_1 & a_{12}A_2 & \cdots & a_{1n}A_n \\ a_{21}A_1 & a_{22}A_2 & \cdots & a_{2n}A_n \\ \vdots & \vdots & \vdots & \vdots \\ a_{n1}A_1 & a_{n2}A_2 & \cdots & a_{nn}A_n \end{bmatrix}$$

Proof. A direct calculation gives that $B^*B = I$. \qed

Remark 5.3. The above theorem can be generalized to tight frames - which we do in a later section.

5.3. Third Case. In this section we find unitary matrices of the form

$$\begin{bmatrix} a & a & a & -b & -b & b \\ a & a & a & -b & b & -b \\ a & a & a & b & -b & -b \\ -b & b & b & a & a & a \\ b & -b & b & a & a & a \\ b & b & -b & a & a & a \end{bmatrix}$$

For this to work, we need:

$$3a^2 + 3b^2 = 1,$$

and for the inner products to be zero,

$$3a^2 - b^2 = 0.$$

Solving we get:

$$b = \pm \frac{1}{2}, \text{ and } a = \pm \sqrt{\frac{1}{12}}.$$
6. Two Tight Frames

Here we will give a general construction for unit norm 2-tight frames. The FRC needed such examples made up of a single vector. Unit norm 2-tight frames are the row vectors of a $2n \times n$ matrix having the following properties.

1. The rows square sum to one.
2. The columns square sum to two.
3. The column vectors are orthogonal.

We start with a concrete example and then go to the general case.

\[
\begin{bmatrix}
-\sqrt{\frac{7}{16}} & \sqrt{\frac{5}{16}} & \sqrt{\frac{3}{16}} & \sqrt{\frac{1}{16}} \\
\sqrt{\frac{7}{16}} & -\sqrt{\frac{5}{16}} & \sqrt{\frac{3}{16}} & \sqrt{\frac{1}{16}} \\
\sqrt{\frac{7}{16}} & \sqrt{\frac{5}{16}} & -\sqrt{\frac{3}{16}} & \sqrt{\frac{1}{16}} \\
\sqrt{\frac{7}{16}} & \sqrt{\frac{5}{16}} & \sqrt{\frac{3}{16}} & -\sqrt{\frac{1}{16}} \\
-\sqrt{\frac{1}{16}} & \sqrt{\frac{3}{16}} & \sqrt{\frac{5}{16}} & \sqrt{\frac{7}{16}} \\
\sqrt{\frac{1}{16}} & -\sqrt{\frac{3}{16}} & \sqrt{\frac{5}{16}} & \sqrt{\frac{7}{16}} \\
\sqrt{\frac{1}{16}} & \sqrt{\frac{3}{16}} & -\sqrt{\frac{5}{16}} & \sqrt{\frac{7}{16}} \\
\sqrt{\frac{1}{16}} & \sqrt{\frac{3}{16}} & \sqrt{\frac{5}{16}} & -\sqrt{\frac{7}{16}}
\end{bmatrix}
\]

If we let

\[
A = \begin{bmatrix}
-\sqrt{\frac{7}{16}} & \sqrt{\frac{5}{16}} & \sqrt{\frac{3}{16}} & \sqrt{\frac{1}{16}} \\
\sqrt{\frac{7}{16}} & -\sqrt{\frac{5}{16}} & \sqrt{\frac{3}{16}} & \sqrt{\frac{1}{16}} \\
\sqrt{\frac{7}{16}} & \sqrt{\frac{5}{16}} & -\sqrt{\frac{3}{16}} & \sqrt{\frac{1}{16}} \\
\sqrt{\frac{7}{16}} & \sqrt{\frac{5}{16}} & \sqrt{\frac{3}{16}} & -\sqrt{\frac{1}{16}} \\
-\sqrt{\frac{1}{16}} & \sqrt{\frac{3}{16}} & \sqrt{\frac{5}{16}} & \sqrt{\frac{7}{16}} \\
\sqrt{\frac{1}{16}} & -\sqrt{\frac{3}{16}} & \sqrt{\frac{5}{16}} & \sqrt{\frac{7}{16}} \\
\sqrt{\frac{1}{16}} & \sqrt{\frac{3}{16}} & -\sqrt{\frac{5}{16}} & \sqrt{\frac{7}{16}} \\
\sqrt{\frac{1}{16}} & \sqrt{\frac{3}{16}} & \sqrt{\frac{5}{16}} & -\sqrt{\frac{7}{16}}
\end{bmatrix}
\]

We can think of reversing the rows of $A$ to get the matrix:

\[
A^R = \frac{1}{\sqrt{16}} \begin{bmatrix}
\sqrt{7} & \sqrt{5} & \sqrt{3} & -\sqrt{1} \\
\sqrt{7} & \sqrt{5} & -\sqrt{3} & \sqrt{1} \\
\sqrt{7} & -\sqrt{5} & \sqrt{3} & \sqrt{1} \\
-\sqrt{7} & \sqrt{5} & \sqrt{3} & \sqrt{1}
\end{bmatrix}
\]

We can also think of reversing the columns of $A$ to get:

\[
A^C = \frac{1}{16} \begin{bmatrix}
\sqrt{7} & \sqrt{3} & \sqrt{3} & -\sqrt{7} \\
\sqrt{7} & \sqrt{3} & -\sqrt{5} & \sqrt{7} \\
\sqrt{7} & -\sqrt{3} & \sqrt{5} & \sqrt{7} \\
-\sqrt{7} & \sqrt{3} & \sqrt{5} & \sqrt{7}
\end{bmatrix}
\]
If we reverse the rows of $A$ and then reverse the columns of the resulting matrix we get:

$$A^{RC} = \begin{bmatrix}
\sqrt{\frac{1}{16}} & \sqrt{\frac{3}{16}} & \sqrt{\frac{5}{16}} & \sqrt{\frac{7}{16}} \\
\frac{1}{16} & -\sqrt{\frac{3}{16}} & \sqrt{\frac{5}{16}} & \sqrt{\frac{7}{16}} \\
\frac{1}{16} & \sqrt{\frac{3}{16}} & -\sqrt{\frac{5}{16}} & \sqrt{\frac{7}{16}} \\
\frac{1}{16} & \sqrt{\frac{3}{16}} & \sqrt{\frac{5}{16}} & -\sqrt{\frac{7}{16}} \\
\frac{1}{16} & -\sqrt{\frac{3}{16}} & \sqrt{\frac{5}{16}} & -\sqrt{\frac{7}{16}} \\
\frac{1}{16} & \sqrt{\frac{3}{16}} & -\sqrt{\frac{5}{16}} & -\sqrt{\frac{7}{16}} \\
\frac{1}{16} & -\sqrt{\frac{3}{16}} & \sqrt{\frac{5}{16}} & -\sqrt{\frac{7}{16}} \\
\frac{1}{16} & \sqrt{\frac{3}{16}} & -\sqrt{\frac{5}{16}} & -\sqrt{\frac{7}{16}}
\end{bmatrix}$$

So our first matrix is of the form:

$$\begin{bmatrix} A \\ A^{RC} \end{bmatrix}$$

Let us do this one more time.

$$\frac{1}{\sqrt{64}} \begin{bmatrix} -\sqrt{15} & \sqrt{13} & \sqrt{11} & \sqrt{9} & -\sqrt{7} & \sqrt{5} & \sqrt{3} & \sqrt{1} \\
\sqrt{15} & -\sqrt{13} & \sqrt{11} & \sqrt{9} & \sqrt{7} & -\sqrt{5} & \sqrt{3} & \sqrt{1} \\
\sqrt{15} & \sqrt{13} & -\sqrt{11} & \sqrt{9} & \sqrt{7} & \sqrt{5} & -\sqrt{3} & \sqrt{1} \\
\sqrt{15} & -\sqrt{13} & -\sqrt{11} & \sqrt{9} & \sqrt{7} & -\sqrt{5} & -\sqrt{3} & -\sqrt{1} \\
-\sqrt{15} & \sqrt{13} & \sqrt{11} & \sqrt{9} & \sqrt{7} & \sqrt{5} & \sqrt{3} & -\sqrt{1} \\
\sqrt{15} & -\sqrt{13} & \sqrt{11} & \sqrt{9} & -\sqrt{7} & \sqrt{5} & -\sqrt{3} & -\sqrt{1} \\
\sqrt{15} & \sqrt{13} & -\sqrt{11} & \sqrt{9} & -\sqrt{7} & -\sqrt{5} & \sqrt{3} & -\sqrt{1} \\
\sqrt{15} & -\sqrt{13} & -\sqrt{11} & \sqrt{9} & -\sqrt{7} & -\sqrt{5} & -\sqrt{3} & \sqrt{1}
\end{bmatrix}$$

In our earlier notation, this is of the form:

$$\begin{bmatrix} A & B \\ A & -B \\ B^{RC} & A^{RC} \\ B^{RC} & -A^{RC} \end{bmatrix}$$

This construction will always work as long as the top row is an arithmetic progression. So the general case here looks like (in the $8 \times 4$ case):
The columns are orthogonal, the rows square sum to 1 and the columns square sum to 2. i.e. This is a 2-tight frame. We can then iterate this procedure as above to get ever larger examples.

7. TWO TIGHT FRAMES: WEIGHT IN FRONT

The idea now is to create 2-tight frames which have a big portion of their weight at the beginning. A group at the FRC wanted this for their work on the Kadison-Singer Problem. This will give an alternative concrete example of non-2-pavable projections. In [1] it was first shown that the class of projections with constant diagonal 1/2 are not two-pavable. Later, in [2] a concrete construction of such projections was given using variations of the discrete Fourier transform.

Let us start with an example and then see that it is a very general concept.

\[
\frac{1}{\sqrt{64}}\begin{bmatrix}
-\sqrt{15} & \sqrt{14} & \sqrt{13} & \sqrt{12} & -\sqrt{4} & \sqrt{3} & \sqrt{2} & \sqrt{1} \\
\sqrt{15} & -\sqrt{14} & \sqrt{13} & \sqrt{12} & \sqrt{4} & -\sqrt{3} & \sqrt{2} & \sqrt{1} \\
\sqrt{15} & \sqrt{14} & -\sqrt{13} & \sqrt{12} & \sqrt{4} & \sqrt{3} & -\sqrt{2} & \sqrt{1} \\
-\sqrt{15} & \sqrt{14} & \sqrt{13} & -\sqrt{12} & \sqrt{4} & \sqrt{3} & \sqrt{2} & -\sqrt{1} \\
\sqrt{15} & -\sqrt{14} & \sqrt{13} & \sqrt{12} & -\sqrt{4} & -\sqrt{3} & -\sqrt{2} & -\sqrt{1} \\
\sqrt{15} & \sqrt{14} & -\sqrt{13} & \sqrt{12} & -\sqrt{4} & -\sqrt{3} & -\sqrt{2} & -\sqrt{1} \\
-\sqrt{15} & -\sqrt{14} & \sqrt{13} & -\sqrt{12} & -\sqrt{4} & \sqrt{3} & -\sqrt{2} & \sqrt{1} \\
\sqrt{15} & \sqrt{14} & \sqrt{13} & \sqrt{12} & -\sqrt{4} & -\sqrt{3} & -\sqrt{2} & \sqrt{1}
\end{bmatrix}
\]

We note that in the matrix above, the columns are orthogonal and the columns square sum to 2 and the rows square sum to 1.

Now we will seriously generalize this example by understanding why it is working. What makes this work comes from the proof of how we get the
formula for $1 + 2 + 3 + \cdots + n$. To do this, we write these numbers in reverse and add them in pairs

$$
\begin{bmatrix}
1 & 2 & 3 & \cdots & n - 2 & n - 1 & n \\
n & n - 1 & n - 2 & \cdots & 3 & 2 & 1
\end{bmatrix}
$$

We note that the sum of the columns are all the same and they add up to twice the sum of the rows. So we will work with integers $n_1 > n_2 > \ldots > n_{2k-1}$ and decide what we want the sums to equal, say $m$, and consider the matrix

$$
\begin{bmatrix}
\sqrt{n_1} & \sqrt{n_2} & \cdots & \sqrt{n_{2k-1}} & \sqrt{m-n_{2k-1}} & \sqrt{m-n_{2k-1-1}} & \cdots & \sqrt{m-n_1}
\end{bmatrix}
$$

Then we write the first row $2^k$ times and below it the second row $2^k$ times. We now need to assign signs of $\pm$ to this $2^{k+1} \times 2^k$ matrix so that the columns are orthogonal. This is easy to do. We use a variation of Hadamard matrices. i.e.

$$
A_1 = \begin{bmatrix}
-1 & 1 & 1 & 1 \\
1 & -1 & 1 & 1 \\
1 & 1 & -1 & 1 \\
1 & 1 & 1 & -1
\end{bmatrix}
$$

And by induction,

$$
A_{n+1} = \begin{bmatrix}
A_n & A_n \\
A_n & -A_n
\end{bmatrix}
$$

We use this twice, i.e. for each set of $2^k$ copies of the two rows. With this class of signs, we get a $2^{k+1} \times 2^k$ matrix with orthogonal columns and square row sums

$$
(1) \quad \sum_{i=1}^{2^{k-1}} n_i + \sum_{i=1}^{2^{k-1}} (m - n_i) = 2^{k-1}m,
$$

and square sums of the columns is just $2^k m$. So after normalization, this is a unit norm 2-tight frame.

Our examples above do this. Also, the following works - for example:

$$
\begin{array}{cccccccc}
32 & 31 & 30 & 29 & 25 & 24 & 23 & 22 \\
18 & 17 & 16 & 15 & 11 & 10 & 9 & 8 \\
8 & 9 & 10 & 11 & 15 & 16 & 17 & 18 \\
22 & 23 & 24 & 25 & 29 & 30 & 31 & 32
\end{array}
$$

Further generalizations:

Remark 7.1. Note that the terms may not all be decreasing above. This would require that

$$
n_{2k-1} > m - n_{2k-1}.
$$
That is,

\[ m < 2n_{2^{k-1}}. \]

Remark 7.2. Note that we do not really need the terms to be decreasing above. All we really need is for the rows to sum to \(2^{k-1}m\) and the columns to sum to \(m\). For example,

\[
\begin{bmatrix}
20 & 24 & 19 & 25 & 1 & 7 & 2 & 6 \\
6 & 2 & 7 & 1 & 25 & 14 & 24 & 20
\end{bmatrix}
\]

Remark 7.3. Note the we do not need the numbers to be distinct above. We can write one of them as many times as we like as long as we satisfy Equation 1.

Remark 7.4. Note that this gives another solution to the paper of Casazza, Fickus, Mixon and Tremain \[2\] where they used variations of the DFT to construct non-2-pavable matrices. A careful choice of the numbers above putting most of the weight in the first \(n-1\) positions then works in their proof.

Here is another outline for constructing 2-tight frames. Fix \(a, b, c, d\) and let \(m = a + b + c + d\).

\[
\begin{bmatrix}
a & b & c & d & m - d & m - c & m - b & m - a \\
1 & 2 & 3 & 4 & 5 & 6 & 7 & 8
\end{bmatrix}
\]

Then, the sum of the column numbers is \(a + b + c + d\) and we have 8 of these so the total sum after we turn this into a matrix is \(8m\) while the sum of the rows is \(4m\). So, after normalization, this is a unit norm 2-tight frame.

8. General Tight Frames

We will generalize one of our earlier examples.

**Theorem 8.1.** Let \(\{A_i\}_{i=1}^n\) be \(r \times k_i\) matrices with orthogonal columns. Let \((a_{ij})_{i=1,j=1}^{m,n}\) be a matrix with orthogonal columns.

Then

\[
\begin{bmatrix}
a_{11}A_1 & a_{12}A_2 & \cdots & a_{1n}A_n \\
a_{21}A_1 & a_{22}A_2 & \cdots & a_{2n}A_n \\
\vdots & \vdots & \ddots & \vdots \\
a_{m1}A_1 & a_{m2}A_2 & \cdots & a_{mn}A_n
\end{bmatrix}
\]

is a \(mr \times L\) matrix with orthogonal columns where

\[ L = \sum_{i=1}^{n} k_i. \]

**Proof.** Obvious. \(\square\)

This leads to:
Theorem 8.2. Let $A_1, A_2, \ldots, A_n$ be $km \times m$ matrices with unit norm rows and orthogonal columns which square sum to $km$. i.e. These are unit norm $k$-tight frames. Let 

$$(a_{ij})_{i,j=1}^n,$$

be an orthonormal matrix.

Then the following matrix is a $kmn \times mn$ unit norm tight frame:

$$
\begin{pmatrix}
a_{11}A_1 & a_{12}A_2 & \cdots & a_{1n}A_n \\
a_{21}A_1 & a_{22}A_2 & \cdots & a_{2n}A_n \\
\vdots & \vdots & \ddots & \vdots \\
a_{n1}A_1 & a_{n2}A_2 & \cdots & a_{nn}A_n
\end{pmatrix}
$$

Proof. Obvious. 

□

Remark 8.3. The same theorems hold with rows replaced by columns.

9. Unbiased Bases

Here is a way to get three unbiased bases in $\mathbb{R}^4$. One will be the unit vectors. The other two are below where $+, -$ means $±1$:

$$
\begin{pmatrix}
- & + & + & + \\
+ & - & + & + \\
+ & + & - & + \\
+ & + & + & -
\end{pmatrix}
$$

And the last basis:

$$
\begin{pmatrix}
+ & + & + & + \\
+ & + & - & - \\
+ & - & - & - \\
+ & - & + & -
\end{pmatrix}
$$

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Department of Mathematics, University of Missouri, Columbia, MO 65211-4100

E-mail address: j.tremain@mchsi.com