Overview

Computer Supported Query Formulation

Asymetrix Report 94-5

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Version of May 20, 2021 at 0:44

PUBLISHED AS:

H.A. (Erik) Proper. An Overview of Computer Supported Query Formulation. Technical report, Asymetrix Research Laboratory, University of Queensland, Brisbane, Queensland, Australia, 1994.

1 Introduction

Most present day organisations make use of some automated information system. This usually means that a large body of vital corporate information is stored in these information systems. As a result, an essential function of information systems should be the support of disclosure of this information.

We purposely use the term information disclosure in this context. When using the term information disclosure we envision a computer supported mechanism that allows for an easy and intuitive formulation of queries in a language that is as close to the user’s perception of the universe of discourse as possible. From this point of view, it is only obvious that we do not consider a simple query mechanism where users have to enter complex queries manually and look up what information is stored in a set of relational tables. Without a set of adequate information disclosure avenues an
information system becomes worthless since there is no use in storing information that will never be retrieved.

An adequate support for information disclosure, however, is far from a trivial problem. Most query languages and query mechanisms do not provide any support for the users in their quest for information. Most of these existing mechanisms can hardly be called disclosure mechanisms as they do not provide users any support during the formulation process. Furthermore, the conceptual schemata of real-life applications tend to be quite large and complicated. As a result, users may easily become lost in conceptual space and they will end up retrieving irrelevant (or even wrong) objects and may miss out on relevant objects. Retrieving irrelevant objects leads to a low precision, missing relevant objects has a negative impact on the recall ([SM83]).

The disclosure of information stored in an information system has some clear parallels to the disclosure problems encountered in document retrieval systems. To draw this parallel in more detail, we quote the information retrieval paradigm as introduced in [BW92]. The paradigm starts with an individual or company having an information need they wish to fulfil. This need is typically a vague notion and needs to be made more concrete in terms of an information request (the query) in some (formal) language. The information request should be as good as possible a description of the information need. The information request is then passed on to an automated system, or a human intermediary, who will then try to fulfil the information request using the information stored in the system. This is illustrated in the information disclosure, or information retrieval paradigm, presented in figure 1 which is taken from [BW92].

We now briefly discuss why the information retrieval paradigm for document retrieval systems is also applicable for information systems. For a more elaborate discussion on the relation between information systems and document (information) retrieval systems in the context of the information retrieval paradigm, refer to [Pro94a]. In the paradigm, the retrievable information is modelled as a set \( K \) of information objects constituting the information base (or population).

In a (multi-media) document retrieval system the information base will be a set of multi-media documents ([SM83]), while in the case of an information system the information base will contain a set of facts conforming to a conceptual schema (although this could be multi-media as well). Each information object \( o \in K \) is characterised by a set of descriptors \( \mathcal{X}(o) \) that facilitates its disclosure. The characterisation of information objects is carried out by a process referred to as indexing. In an information system, the stored objects (the population or information base) can always be identified by a set of (denotable) values, the identification of the object. For example, an address may be identified as a city name, street name, and house number. The characterisation of objects in an information system is directly provided by the reference schemes of the object types.

The actual information disclosure is driven by a process referred to as matching. In document retrieval applications this matching process tends to be rather complex. The
characterisation of documents is known to be a hard problem ([Mar77], [Cra86]), although newly developed approaches turn out to be quite successful ([Sal89]). In information systems the matching process is less complex as the objects in the information base have a more clear characterisation (the identification). In this case, the identification of the objects (facts) is simply related to the query formulation $q$ by some (formal) query language.

The remaining problem is the query formulation process itself. An easy and intuitive way to formulate queries is absolutely essential for an adequate information disclosure. Quite often, the quest from users to fulfil their information need can be aptly described by ([Bru93]):

*I don’t know what I’m looking for, but I’ll know when I find it.*

In document retrieval systems this problem is attacked by using *query by navigation* ([BW92], [Bru93]) and *relevance feedback* mechanisms ([Rij89]). The query by navigation interaction mechanism between a searcher and the system is well-known from the Information Retrieval field, and has proven to be useful. It shall come as no surprise that these mechanisms also apply to the query formulation problem for information systems. In [BPW93], [BPW94], [HPW94b], [Pro94a] such applications of the *query by navigation* and *relevance feedback* mechanisms have been described before. When combining the query by navigation and manipulation mechanisms with the ideas behind visual interfaces for query formulation as described in e.g. [ADD92] and [Ros94], powerful and intuitive tools for computer supported query formulation become feasible, resulting in improved information disclosure. Such tools will also
heavily rely on the ideas of direct manipulation interfaces ([Sch83]) as used in present
day computer interfaces.

One important step in the improvement of the information disclosure of information
systems, is the introduction of query languages on a conceptual level. These lan-
guages allow for the formulation of queries in terms common to the users, i.e. the
verbalisations of the types in the conceptual schema. Examples of such conceptual
query languages are RIDL ([Mee82]), LISA-D ([HPW93], [HPW94a]), and FORML
([HHO92]). By letting users formulate queries on a conceptual level, users are safe-
guarded from having to know the exact mapping to internal representations (e.g. a set
of tables which conform to the relational model) to be able to formulate queries in a
non conceptual language such as SQL. The next step is to introduce ways to support
users in the formulation of queries in such conceptual query languages (CQL).

2 A New Generation of Formulation Mechanisms

In line with the above discussed information retrieval paradigm and the notion of rel-
evance feedback, a query formulation process (both for a document retrieval system,
and an information system) can be said to roughly consist of the following four phases:

1. The explorative phase. What information is there, and what does it mean?

2. The constructive phase. Using the results of phase 1, the actual query is formu-
lated.

3. The feedback phase. The result from the query formulated in phase 2 may not
be completely satisfactory. In this case, phases 1 and 2 need to be re-done and
the result refined.

4. The presentation phase. In most cases, the result of a query needs to be incorpo-
rated into a report or some other document. This means that the results must be
grouped or aggregated in some form.

Depending on the user’s knowledge of the system, the importance of the respective
phases may change. For instance, a user who has a good working knowledge of the
structure of the stored information may not require an elaborate first phase and would
like to proceed with the second phase as soon as possible.

In the research for the InfoAssistant product, we try to integrate a palette of comple-
mentary mechanisms to formulate queries on a conceptual level. These mechanisms
are the following:

**query by navigation** This mechanism has been introduced in [BPW93], [PW95] and
[Pro94a]. The idea behind this mechanism is to shape a conceptual schema,
which is essentially a graph, as a hypertext and letting users formulate (part of) their information need by navigating through this hypertext. This mechanisms is particularly suited for those users who do not have a clear idea of what information is stored in the information system as it is able to truly guide the user through the (structure of the) stored information.

A precursor of the query by navigation mechanism for information systems exists for information retrieval systems ([Bru93]). In experiments it was shown that in the IR case, this mechanism helps novice users in finding their way around the stored information, without hampering expert users ([BBB91]).

All research that remains to be done in this area is some tuning and adapting the existing (academic) ideas to the applied situation in InfoAssistant.

**query by construction** This mechanism has also been discussed before in [BPW93], [PW95] and [Pro94a]. This mechanism was born out of the observation that the results of a query by navigation sessions are relatively simple queries without advanced operations such as grouping, intersections, counting, etc. Extending the query by navigation mechanisms with such operations would have led to an unacceptable increase in complexity. Therefore the introduction of the query by construction as an additional mechanism was chosen.

The query by construction mechanism is basically a syntax directed editor which allows a user to combine the query particles resulting from query by navigation (and the three additional mechanisms discussed below) sessions to be combined into complex queries using the more advanced operations.

Research-wise, this part is finished as there is not much research needed for a syntax directed editor

**point to point queries** The point to point queries originated from a rough idea from J. Harding. A point to point query starts by selecting two or more object types from a conceptual schema. Then the system should return a list of possible (non cyclic) paths through the information structure between the specified object types. For obvious reasons, the paths in this list should be ordered according to some relevance criterion.

This style of querying corresponds to a situation in which users know some aspects (object types) about which they want to be informed, but do not yet know the exact details of their information need and the underlying information structure. The query by navigation mechanism, on the other hand, is intended to support users who do not have an overview of the stored information.

In [Pro94b] this mechanism is discussed and formalised in full detail.

**spider queries** This mechanism originated from a discussion with L. Delano. Users quite often simply want to know all information about instances of an object type \(x\). For this purpose the spider queries were introduced. A crucial aspect of spider queries is of course limiting the all information as users probably do not
want to be confronted with a listing of all information stored in the information system.

The idea behind spider queries is to start out from one object type, and to associate all information that is relevant to this object type. The essential part of a spider query is selecting the object types in the direct surroundings of the initial object type that are considered to be relevant, thus limiting the amount of information returned to the user.

This style of querying corresponds to a situation where users only know about the existence of some object types in the conceptual schema about which they would like to be informed.

A complete discussion and formal treatment of this mechanism can be found in [Pro94b].

natural language queries A more commonly known mechanism for computer supported query formulation are (semi) natural language query formulation systems. These mechanisms try to interpret sentences in a semi-natural language format and generate an appropriate query in SQL.

Our aim is to try and integrate these ideas with the newly added formulation mechanisms. One important aspect of this integration is that it would allow us to interpret the natural language sentence, and then automatically formulate a query in a conceptual query language rather than SQL. This would certainly put the user in a much better position to validate the resulting query than to confront users with an SQL query.

A natural language formulation mechanism is useful for those users who know what information is stored in the information system, but who do not know the exact names of the types. The flexibility of a semi-natural language would then cater for this.

In the remainder of this overview report we discuss some example session using the different disclosure avenues. This should give a more hands-on idea of what these mechanisms are about.

3 An Example Session

In this section we discuss a sample session using the query formulation component of InfoAssistant. The discussed example operates on a conceptual schema for the administration of the election of American presidents. The example schema itself is not shown; the structure of the domain will become clear from the sample session. Note that the quality of the verbalisations of paths expressions used in the examples in this section should be improved. However, this is the subject of further research.
In figure 2, a possible screen is depicted for building queries using a point to point query mechanism. The upper window is concerned with the point to point query itself, whereas the lower window contains the complete query under construction. When specifying a point to point query a user specifies a sequence of object types: the points. For each point, the user is offered a listbox containing all object types present in the conceptual schema. The order of the object types in the listbox should preferably be based on some notion of conceptual importance ([CH94]). In figure 3 an existing point to point query path from president to election is extended with another point.

After all points of the point to point query have been specified, the point to point query can be transformed into a proper query (i.e. a path through the conceptual schema) by pressing the Go! button in the point to point query window. In figure 4, this process is illustrated. The sample PPQ involves three points. Therefore, two paths through the conceptual schema will result. We now shift our attention from the point to point query window to the query by construction window. Note that the small box containing the PPQ abbreviation is now replaced by the paths resulting from the point to point query (i.e. President winning election which resulted in nr of votes). The system initially inserts a most likely path. The user can, however, select alternative paths using a listbox. Note that not all alternative paths between the two points are listed in the listbox. The reason for this is the NP completeness of the path searching problem. To avoid the NP completeness problem, only the best paths are listed initially. However, potentially all paths can be selected (which still remains NP complete) by repeatedly selecting the MORE option. In the remainder of this article we will discuss this in more detail.

Since every path resulting from a query by navigation session connects two points in the conceptual schema, any path through the conceptual schema displayed in the query by construction screen can be used as a starting point for a query by navigation session, and vice versa. This is illustrated in figure 5. In this session, the user has selected the box which contains the two paths politician is president of administration and inaugurated in year for a query by navigation session. The upper window now displays a node in the query by navigation session, with the path politician is president of administration inaugurated in year as its focus. If the user had selected the inaugurated in year listbox, the initial focus would have been administration inaugurated in year.

The query by construction window in figure 5 basically offers a syntax directed editor. In the left part of the window all possible constructs from the query language are listed. In our examples we have used the constructs defined in LISA-D. Once the FORML and LISA-D languages have been merged, a more complete language for the query by construction part will result.

Next we discuss a session involving a spider query. We start out from an existing query in a query by construction window, which could have been constructed using a query by navigation query or a point to point query. Note that this could also be single object type, e.g. politician. The spider query mechanism adds one important aspect to the query by construction window, the spider button: [...]. When a user presses this button, the system calculates the spider query of the object type directly to the right of
the button. This is illustrated in figure 7. The system allows for the removal of parts of
the resulting spider query that are not considered to be relevant by the user. Suppose
the user is not interested in administration is headed by and election won by, then these
paths can be deleted, which leads to the screen depicted in figure 8.

It is now interesting to see that a query essentially is a double tree with a shared root
(politician in the example). Furthermore, the leaves on the tree resulting from the spider
query can be extended further if desired by commencing new spider queries. Finally,
since the result of a spider query is constructed from path expressions as well, these
expressions have the $\mathbb{R}$ associated that can be used to select alternative paths between
the head and tail object types. Furthermore, the paths can also be used as a starting
point of a query by navigation session. This latter possibility is illustrated in figure 9.

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Figure 2: Building a PPQ query
Figure 3: Extending the PPQ path
Figure 4: Completing a PPQ
Figure 5: Switching to query by navigation
Figure 6: Start of a spider query
Figure 7: Result of a Spider Query
Figure 8: Pruning the Spider Query
Focus: the politician which served the nr of years

the politician
the nr of years
the politician which served the nr of years that were served by the president
the president which served the nr of years
nr of years that were served by the politician

AND ALSO

inaugurated in year

1920

Go!

the politician which served the nr of years

nr of years that were served by the politician

win the election

which resulted in nr of votes

is president of administration

Query by Navigation

Figure 9: Switching to query by navigation