Symmetry reduction in the augmented space recursion formalism for random binary alloys

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Abstract. We present here an efficient method which systematically reduces the rank of the augmented space and thereby helps to implement augmented space recursion for any real calculation. Our method is based on the symmetry of the Hamiltonian in the augmented space and keeping recursion basis vectors in the irreducible subspace of the Hilbert space.

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1. Introduction

In a series of earlier communications [1, 2] we have proposed the augmented space recursion (ASR) as an efficient computational technique for the calculation of configuration averaged electronic properties of random binary alloys. The method is a combination of the augmented space formalism [3] and the recursion method of Haydock et al [4]. When coupled with the local density functional approximation based tight-binding linearized muffin-tin orbital [5, 6] it provides a first-principles determination of electronic structure and total energy calculations for disordered alloys. Effects such as short-ranged ordering and local lattice distortions due to size mismatch of the constituent atoms can very easily be incorporated in the methodology [7, 8, 9]. Recently the method has been applied to study dispersion and line widths for phonons in disordered alloys. Since recursion can take into account disorder in beyond nearest-neighbour force constants as well as large environmental effects, this application was a step forward in the theory of phonons in disordered alloys [10]. It has been pointed out

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repeatedly in the above publications that the ASR with its terminator approximation\cite{11,12} goes beyond the standard single-site mean field theories and accurately takes off-diagonal disorder into account. A number of attempts have been made for practical implementation of mean-field theories beyond the single site approximation. One of the most sophisticated developments among the mean-field theories, is the travelling cluster approximation of Gray and Kaplan\cite{13,14} and its latest version: the itinerant coherent potential approximation (ICPA)\cite{15}, where a small number of classes of selected excitations in augmented space are preserved. In spite of the simplification so introduced, the resulting equations are still formidable enough to discourage further generalization to larger clusters beyond the nearest neighbour. In the cluster CPA proposed by Mookerjee\cite{16,17}, in which one retains disorder fluctuations only from a chosen cluster and replaces the rest of the system by an effective medium, one is restricted by the limitation of solving number of coupled self-consistent equations. The number of equations increases with the size of chosen cluster, which again discouraged extension to larger clusters beyond the nearest neighbour. This brings us back to the attempt at the implementation of the ASR for large environment fluctuations, implying accurate evaluation of a large sequence of recursion coefficients. We shall discuss this in some detail in the following section.

The ASR, in its practical implementation to real alloy systems, has several practical problems. The formalism maps a disordered Hamiltonian, described in a Hilbert space spanned by a countable basis in which preferably its representation is sparse, onto an ordered Hamiltonian in an enlarged space. This enlarged or augmented space is constructed by taking the outer product of the configuration space of the random variables of the disordered Hamiltonian together with original Hilbert space in which the disordered Hamiltonian was described. The main difficulty in the implementation of recursion in this augmented space is its enormous rank. A system of $N$ sites with binary distribution has an augmented space of rank $N \times 2^N$. This enormous rank of the augmented space and its description on the computer is what had discouraged earlier workers in implementing the ASR, although the augmented space formalism was first proposed about thirty years ago\cite{3}. The second problem was the recursive propagation of errors. Very small errors during recursion (inevitable in numerical implementations) tends to add up leading to loss of orthogonality in the recursion basis. This leads to appearance of ghost states. Haydock and Te\cite{18} has overcome this difficulty in proposing the dynamical recursion. We refer the reader to their paper for details.

In this communication we shall describe the reduction of the effective rank of the augmented space on which the recursion is carried out, using the local point-group symmetries in both the Hilbert space and the configuration space. Such reduction is well known in various manipulations in reciprocal space for ordered systems. However, such a reduction has not been implemented in real space recursion, although it had been proposed by Gallagher\cite{19} quite some time ago. It had been shown by Gallagher that if we start recursion with a state belonging to an irreducible subspace of a Hilbert
space, subsequent recursion always stays within this subspace. This allows us significant reduction in rank of the required subspace and makes this method practically feasible. Further, this method with its terminator approximation, retains the Herglotz properties of the configuration averaged Green function.

2. The augmented space recursion

2.1. The recursion method

The recursion method addresses inversions of infinite matrices [4]. Once a sparse representation of an operator in Hilbert space, $\tilde{H}$, is known in a countable basis, the recursion method obtains an alternative basis in which the operator becomes tridiagonal. This basis and the representations of the operator in it are found recursively starting from $|u_1\rangle = |R\rangle \otimes |1\rangle$:

\[
|u_2\rangle = \tilde{H}|u_1\rangle - \alpha_1|u_1\rangle
\]

\[
|u_{n+1}\rangle = \tilde{H}|u_n\rangle - \alpha_n|u_n\rangle - \beta^2_n|u_{n-1}\rangle \quad n \geq 2
\]

where,

\[
\alpha_n = \frac{\langle u_n | \tilde{H} | u_n \rangle}{\langle u_n | u_n \rangle}, \quad \beta^2_n = \frac{\langle u_n | u_n \rangle}{\langle u_{n-1} | u_{n-1} \rangle}
\]

for $n \geq 2$. In this basis the representation of $\tilde{H}$ is tridiagonal:

\[
\frac{\langle u_n | \tilde{H} | u_n \rangle}{\langle u_n | u_n \rangle} = \alpha_n \quad ; \quad \frac{\langle u_{n-1} | \tilde{H} | u_n \rangle}{\langle u_{n-1} | u_{n-1} \rangle}^{1/2} = \beta_n \quad ; \quad \langle u_n | \tilde{H} | u_m \rangle = 0 \quad m \geq n + 2
\]

and the averaged Green function can be written as a continued fraction:

\[
\ll G_{RR}(E) \gg = \langle u_1 | (E \tilde{I} - \tilde{H})^{-1} | u_1 \rangle
\]

\[
= \frac{1}{E - \alpha_1 - \frac{\beta^2_2}{E - \alpha_2 - \frac{\beta^2_3}{\ddots}}}
\]

\[
= \frac{1}{E - \alpha_N - T(E)}
\]

(3)
The asymptotic part of the continued fraction is obtained from the initial set of coefficients \( \{\alpha_n, \beta_n\} \) for \( n < N \), using the idea of terminators \([11,12]\). For small values of \( N \) we have large inaccuracies and the more the structure in the spectral density of \( \tilde{H} \) the larger is the \( N \) needed to maintain the window of accuracy required by us.

### 2.2. The TB-LMTO Hamiltonian for random binary alloys

In earlier communications we have described how to deal with random binary alloys \([1]\) within the framework of the tight-binding linearized muffin-tin orbital method (TB-LMTO) \([5,6]\) and using the augmented space formalism (ASF). We refer the readers to the seminal papers mentioned earlier for details. Here we shall introduce only the salient features which will be required by us subsequently in this communication. We begin by setting up a muffin-tin potential with centres at the atomic sites \( R \) on a lattice. Next, we inflate the muffin-tins into atomic spheres and start from a most-localized TB-LMTO representation of the Hamiltonian:

\[
H = \sum_{RL} C_{RL} \mathcal{P}_{RL} + \sum_{RL} \sum_{R'L'} \Delta_{RL}^{1/2} S_{RL,R'L'} \Delta_{R'L'}^{1/2} \mathcal{T}_{RL,R'L'}
\]

The potential parameters, \( C \) and \( \Delta \), are diagonal matrices in the angular momentum indeces, and

\[
\begin{align*}
C_{RL} &= C_A^{\ell} n_R + C_B^{\ell} (1 - n_R) \\
\Delta_{RL}^{1/2} &= (\Delta_A^{\ell})^{1/2} n_R + (\Delta_B^{\ell})^{1/2} (1 - n_R)
\end{align*}
\]

Here \( L \) is a composite label \( \{\ell, m, m_s\} \) for the angular momentum quantum numbers. \( n_R \) is the random site-occupation variable which takes values 0 or 1 with probability \( x \) or \( y \) \( (x + y = 1) \) respectively, depending upon whether the muffin-tin labelled by \( R \) is occupied by \( A \) or \( B \)-type of atom.

The \( \mathcal{P}_{RL} \) and \( \mathcal{T}_{RL,R'L'} \) are projection and transfer operators in the Hilbert space \( \mathcal{H} \) spanned by the tight-binding basis \( \{|RL\} \).

### 2.3. The augmented space formalism within the TB-LMTO

We can associate the random variable \( n_R \) with an operator \( \mathcal{M}_R \) whose eigenvalues correspond to the observed values of \( n_R \), and whose corresponding eigenvectors, \( \{|0_R\}, |1_R\} \) are the state vectors of the variable. These state vectors of the set of \( N \) random variables \( n_R \) of rank \( 2^N \) span a space called configuration space \( \Phi \) with configurations of the kind \( |\uparrow\uparrow\downarrow\ldots \downarrow\uparrow\ldots\rangle \) where...
If we define the configuration $| \uparrow \uparrow \ldots \uparrow \ldots \rangle$ as the reference configuration, then any other configuration may be uniquely labelled by the cardinality sequence: $\{R_k\}$, which is the sequence of positions where we have a $\downarrow$ configuration. The cardinality sequence of the reference configuration is the null sequence $\{\emptyset\}$.

The augmented space theorem states that

$$
\ll A(\{n_R\}) \gg = < \{\emptyset\} | \tilde{A} | \{\emptyset\} >
$$

(5)

where,

$$
\tilde{A} = \int \ldots \int A(\{\lambda_R\}) \prod dP(\lambda_R) \in \Phi
$$

$P(\lambda_R)$ is the spectral density of the self-adjoint operator $M_R$, which is such that the probability density of $n_R$ is given by:

$$
p(n_R) = -\frac{1}{\pi} \lim_{\delta \to 0} \Im m \langle \uparrow_R | \left( (n_R + i\delta)I - M_R \right)^{-1} | \uparrow_R \rangle
$$

$M_R$ is an operator in the space of configurations $\psi_R$ of the variable $n_R$. This is of rank 2 and is spanned by the states $\{|\uparrow_R\rangle, |\downarrow_R\rangle\}$.

$$
M_R = x \mathcal{P}_\uparrow^R + y \mathcal{P}_\downarrow^R + \sqrt{xy} \left( \mathcal{T}_{\uparrow\downarrow}^R + \mathcal{T}_{\downarrow\uparrow}^R \right)
$$

(6)

The expended Hamiltonian $\tilde{H}$ in the augmented space is constructed by replacing all the random variables $n_R$ by the corresponding operators $M_R$. It is an operator in the augmented space $\Psi = \mathcal{H} \otimes \Phi$. So the ASF maps a disordered Hamiltonian described in a Hilbert space $\mathcal{H}$ onto an ordered Hamiltonian in an enlarged space $\Psi$, where the space $\Psi$ is constructed by augmenting the configuration space $\Phi = \prod_R \psi_R$ of the random variables of the disordered Hamiltonian together with the Hilbert space $\mathcal{H}$ of the disordered Hamiltonian. The configuration space $\Phi$ is of rank $2^N$ if there are $N$ muffin-tin spheres in the system.

We may rewrite the expression for the configuration average of the Green operator as:

$$
\ll G(E) \gg = < \{\emptyset\} | \tilde{D} \left( \hat{A} + \hat{B} + \hat{C} - \hat{S} \right)^{-1} \tilde{D} | \{\emptyset\} >
$$

(7)

where,

$$
\hat{A} = \sum_{RL} A \left[ (E - C_L) / \Delta_L \right] \mathcal{P}_R \otimes \mathcal{P}_L \otimes I
$$
\[ \tilde{B} = \sum_{RL} B \left( \frac{(E - C_L)}{\Delta_L} \right) \mathcal{P}_R \otimes \mathcal{P}_L \otimes \mathcal{P}^R_\downarrow \]

\[ \tilde{C} = \sum_{RL} C \left( \frac{(E - C_L)}{\Delta_L} \right) \mathcal{P}_R \otimes \mathcal{P}_L \otimes \left\{ \mathcal{T}^R_{\uparrow \downarrow} + \mathcal{T}^R_{\downarrow \uparrow} \right\} \]

\[ \tilde{S} = \sum_{RL R' L'} S_{RL R' L'} \mathcal{T}^R_{LR} \otimes \mathcal{T}^L_{LL} \otimes \mathcal{I} \]

\[ \tilde{D} = \sum_{RL} A \left( \frac{\Delta^{1/2}_L}{2} \right) \mathcal{P}_R \otimes \mathcal{P}_L \otimes \mathcal{I} + \sum_{RL} B \left( \frac{\Delta^{1/2}_L}{2} \right) \mathcal{P}_R \otimes \mathcal{P}_L \otimes \mathcal{P}^R_\downarrow \]

\[ \quad + \sum_{RL} C \left( \frac{\Delta^{1/2}_L}{2} \right) \mathcal{P}_R \otimes \mathcal{P}_L \otimes \left\{ \mathcal{T}^R_{\uparrow \downarrow} + \mathcal{T}^R_{\downarrow \uparrow} \right\} \]

\[ \tilde{A}' = \sum_{RL} \{ A \left( \frac{C_L}{\Delta_L} \right) / A \left( \frac{1}{\Delta_L} \right) \} \mathcal{P}_R \otimes \mathcal{P}_L \otimes \mathcal{I} \]

\[ \tilde{B}' = \sum_{RL} \{ B \left[ \left( \frac{C_L - E}{\Delta_L} \right) / A \left( \frac{1}{\Delta_L} \right) \right] \} \mathcal{P}_R \otimes \mathcal{P}_L \otimes \mathcal{P}^R_\downarrow \]

\[ \tilde{C}' = \sum_{RL} \{ C \left[ \left( \frac{C_L - E}{\Delta_L} \right) / A \left( \frac{1}{\Delta_L} \right) \right] \} \mathcal{P}_R \otimes \mathcal{P}_L \otimes \left\{ \mathcal{T}^R_{\uparrow \downarrow} + \mathcal{T}^R_{\downarrow \uparrow} \right\} \]

\[ \tilde{S}' = \sum_{RL R' L'} \{ \left( \frac{1}{A \left( \Delta_L \right)} \right)^{-1/2} S_{RL R' L'} \left[ A \left( \frac{1}{\Delta_L} \right) \right]^{-1/2} \mathcal{T}^R_{RR} \otimes \mathcal{T}^L_{LL} \otimes \mathcal{I} \]
This equation is now exactly in the form in which the recursion method may now be applied. The computational burden is considerably reduced due to this diagonal formulation, the recursion now becomes energy dependent as is clear from the form of the effective Hamiltonian as shown in (7) and discussed in [21, 22]. The recursion formalism of the ordered Hamiltonian was free of this constraint. This energy dependence makes the recursion technique computationally unsuitable because to obtain the Green functions we have to carry out recursion per energy point of interest. This problem has been tackled using seed recursion technique [22]. The idea is to choose a few seed points across the energy spectrum uniformly, carry out recursion over those points and then interpolate the values of coefficients across the band. In this way one may reduce computation time. For example, if one is interested in an energy spectrum of 200 points, in the bare diagonal formulation recursion has to be carried out at all the 200 points but in the seed recursion technique one needs to perform recursions only at 10-15 points. The whole idea stems from the fact that in most of the cases of binary alloys, it is seen that the recursion coefficients \( a_n \) and \( b_n \) vary quite weakly across the energy spectrum. At this point we note that the above expression for the averaged \( \ll G_{LL}(E) \gg \) is exact.

3. Symmetry reduction of the augmented space rank

We mentioned earlier that recursion on the augmented space is not computationally feasible because of its large rank. For a binary alloy with \( N \) sites and with only \( s \)-orbitals on them, the rank of the augmented space is \( N \times 2^N \). Implementing recursion on this huge space for a sufficient number of steps to ensure accuracy is often not feasible on available computers. However, if we exploit the symmetry both of the underlying lattice in real space and of the configuration space (which arises due to the homogeneity of the disorder and arrangements of atoms on the underlying lattice), the rank of the irreducible part of the augmented space in which the recursion is effectively carried out becomes tractable. The conceptual advantages in ASF include apart from analyticity, translational and underlying lattice symmetries automatically built into the augmented space Hamiltonian. This allows us to involve the idea of utilizing symmetry operations present in both the real and configuration spaces, in the context of recursion method, reducing the rank of Hamiltonian drastically and making the implementation of ASF feasible. Since the augmented space recursion essentially retains all the properties of real space recursion but is described in a much enlarged space, it will be useful to consider symmetry operations in real space recursion first and then to consider those in the full augmented space.

During the recursion, the basis member \( |u_n\rangle \) is generated from the starting state \( |u_1\rangle \) by repeated application of the Hamiltonian. If the starting state belongs to an irreducible sub-space of the Hilbert space then all subsequent states generated from recursion will belong to the same irreducible sub-space. Physically, we may understand this as follows: the recursion states \( |u_n\rangle \) carry information of distant environment of the starting state.
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If the Hamiltonian is nearest-neighbour only, then the state $|u_2\rangle$, which arises by the application of the Hamiltonian on $|u_1\rangle$ is a combination of states in the nearest shell with which $|u_1\rangle$ couple via the Hamiltonian. Similarly, $|u_n\rangle$ is a combination of $n$-th neighbour shell with which $|u_1\rangle$ is coupled via the Hamiltonian. If $R$ is a point group symmetry of the Hamiltonian, then all $n$-th neighbour-shell states which are related to one another through the symmetry operator must have equal coupling to $|u_1\rangle$. Hence, it is useful to consider among the $n$-th neighbour-shell states of which $|u_n\rangle$ is a linear combination, only those belonging to the irreducible subspace and redefine the Hamiltonian operation.

As an example, take a nearest-neighbour $s$-state Anderson model on a square lattice, with a binary distribution of its diagonal elements. We shall label the tight-binding basis by the position of the lattice points in units of the lattice constant: e.g. $|(mn)\rangle$ where $m, n$ are integers $0, \pm 1, \pm 2, \ldots$. The starting state $|(00)\rangle$ belongs to the one dimensional representation of the point group of a square lattice. This state then couples with linear combinations of states on neighbour shells which are symmetric under square rotations:

$$
|0m\rangle = \left( |(0m)\rangle + |(m0)\rangle + |(m0)\rangle + |(0m)\rangle \right)/2, \quad m > 0 \\
|mm\rangle = \left( |(mm)\rangle + |(m\bar{m})\rangle + |(m\bar{m})\rangle + |(m\bar{m})\rangle \right)/2, \quad m > 0 \\
|mn\rangle = \left( |(mn)\rangle + |(\bar{m}n)\rangle + |(m\bar{n})\rangle + |(m\bar{n})\rangle \right)/2 \\
+ |(nm)\rangle + |(\bar{n}m)\rangle + |(n\bar{m})\rangle + |(n\bar{m})\rangle \right)/2(2\sqrt{2})
$$

$m > 1, \quad 0 < n < m \quad (10)$

The figure shows the grouping together of sites with the local symmetry of the square lattice.
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lattice on the first three nearest neighbour shells of the central site. The first and second groups in \( [10] \) are coloured gray in figure 1 and the last group coloured black. The labels on the groups (shown on the left sides of the equations \( [10] \)) are confined only to the upper right quadrant of the lattice.

If we go up to \( N \) shells (for large \( N \)) there are about \( 2N^2 \) states in the diamond shaped nearest-neighbour cluster. However, there are only \( (N^2/4 + N/2) \sim N^2/4 \) states with square symmetry.

So within this reduction we can work only in 1/8-th of the lattice, provided we incorporate proper weights to the states to reproduce the correct matrix elements. If \( |R\rangle \) and \( |R'\rangle \) are two states coupled to each other via the Hamiltonian, and both belong to the same irreducible subspace, and if \( |R_1\rangle, |R_2\rangle, \ldots, |R_{\mathcal{W}_R}\rangle \) are states obtained by operating on \( |R\rangle \) by the symmetry group operations of the real space lattice, then \( \mathcal{W}_R \) is called the weight associated with the state labelled by \( R \). If we wish to retain only the states in the irreducible subspace and throw out the others and yet obtain the same results, we have to redefine the Hamiltonian matrix elements as follows:

\[
\langle R | H | R' \rangle \rightarrow \sqrt{\frac{\mathcal{W}_{R'}}{\mathcal{W}_R}} \langle R | H | R' \rangle
\]  

(11)

In the above prescription, the new irreducible basis introduced by us, reflects only the symmetry of the underlying lattice and holds good for a model system which has s-like orbitals only. But for a real system, the TB-LMTO minimal basis contains members with other symmetries as well. For example, in a cubic lattice with a spd minimal basis, we have basis members with \( s, p, e_g \) and \( t_{2g} \) symmetries. The symmetry of the orbitals is reflected in the two-centered Slater-Koster integrals and this prohibits overlap integrals at certain positions, called symmetric positions with respect to the overlapping orbitals. A few of these symmetric positions for a simple cubic lattice are shown in figure 2. For the sake of clarity we have shown only the projections on \( x - y \) plane. We have indicated the positive and negative parts of the orbital lobes by different shades. From figure 2\( b \) it is easy to argue that since the Hamiltonian is spherically symmetric and the product \( \phi_s(000)\phi_{p_x}(010) \) is positive in the upper right quadrant and negative in the upper left quadrant :

\[
\langle 0, s | H | 2, p_x \rangle = 0.
\]

The same is true for

\[
\langle 0, s | H | 2, p_y \rangle = 0.
\]

Similarly,

\[
\langle 0, s | H | 3, p_z \rangle = 0,
\]

\[
\langle 0, s | H | 3, p_x \rangle = 0.
\]

On the other hand,

\[
\langle 0, s | H | 1, p_x \rangle = -\langle 0, s | H | \bar{1}, p_x \rangle \neq 0.
\]

The above is an illustration, detailed arguments for the orbital symmetry has been discussed by Harrison [20].

The table 1 provides the conditions for obtaining the orbital based symmetry factor \( \beta_{LL'}(R - R') \) for a lattice with cubic symmetry. We shall describe the direction of \( \vec{R} - \vec{R'} \) by the direction cosines \( (\ell mn) \). The table gives the details for the calculation of the symmetry factor \( \beta_{LL'}(\ell mn) \) on a cubic lattice. A look at the figures 2 and 4 shows the reason for introduction of this factor.
| Matrix Element | Slater-Koster Expression | Symmetry Factor | Condition for it to be zero |
|----------------|--------------------------|-----------------|-----------------------------|
| \(t_{xx}\)     | \(t_{s\sigma}\)          |                   |                             |
| \(t_{xy}\)     | \(l t_{sp\sigma}\)       |                   |                             |
| \(t_{xy}\)     | \(m [\sqrt{3} l m t_{sd\sigma} + (1-2l^2) t_{pd\pi}]\) |                   |                             |
| \(t_{xy}\)     | \(\sqrt{3} l mn (t_{p\sigma} - 2 t_{pd\pi})\) |                   |                             |
| \(t_{x,x^2-y^2}\) | \(l [\sqrt{3}/2 (l^2-m^2) t_{p\sigma} + (1-l^2+m^2) t_{pd\pi}]\) |                   |                             |
| \(t_{x,3z^2-r^2}\) | \(n(l^2-m^2)[\sqrt{3}/2 t_{p\sigma} - t_{pd\pi}]\) |                   |                             |
| \(t_{xy,xy}\)  | \(3l^2m^2 t_{dd\sigma} + (l^2+m^2-4l^2m^2) t_{dd\pi} \cdots \) |                   |                             |
| \(t_{xy,yz}\)  | \(l m(3m^2 t_{dd\sigma} + (1-4m^2) t_{dd\pi}) + (m^2-1) t_{dd\delta}\) |                   |                             |
| \(t_{yz,yz}\)  | \(l m(3m^2 t_{dd\sigma} + (1-4m^2) t_{dd\pi}) + (m^2-1) t_{dd\delta}\) |                   |                             |
| \(t_{yz,z^2-r^2}\) | \(\sqrt{3} l m(3/2(l^2-m^2) t_{dd\sigma} + [1+2(l^2-m^2)] t_{dd\pi} \cdots \) |                   |                             |
| \(t_{xy,3z^2-r^2}\) | \(n(n^2-1/2(l^2+m^2)] t_{dd\sigma} - n^2 t_{dd\pi} \cdots \) |                   |                             |
| \(t_{x^2-y^2,y^2}\) | \(3/4(l^2-m^2)^2 t_{dd\sigma} + [l^2+m^2-(l^2-m^2)^2] t_{dd\pi} \cdots \) |                   |                             |
| \(t_{x^2-y^2,3z^2-r^2}\) | \(\sqrt{3}(l^2-m^2)^2 [n^2/2-1/4(l^2+m^2)] t_{dd\sigma} - n^2 t_{dd\pi} \cdots \) |                   |                             |
| \(t_{3z^2-r^2,3z^2-r^2}\) | \(n^2-1/2(l^2+m^2)] t_{dd\sigma} + 3n^2(l^2+m^2) t_{dd\pi} \cdots \) |                   |                             |

**Table 1.** Table showing the Slater-Koster parameters and the consequent symmetry factors \(\beta_{\ell m n}^{\ell m n}\). The parameters for which the zero condition have been omitted can be obtained from the last column by permuting the direction cosines.
Figure 2. $sp$-orbitals on a cubic lattice. Only projections on the $xy$ plane are shown. The positive parts of the lobes are unshaded, while the negative parts of the lobes are shaded. The values of the symmetry factor for the overlaps ($\beta_{LL'}(\ell mn)$, defined in the text) are shown for each direction. The numbering of sites on the cubic lattice is shown in the inset of the right bottom.

For the $s$-orbital overlaps the symmetry factor is always 1 (figure 1(a)). All six neighbours in the cubic lattice in the directions $(100), (\bar{1}00), (010), (0\bar{1}0), (001)$ and $(00\bar{1})$ are equivalent with the same overlap integrals $t_{ss\sigma}$ (only four are shown in the $xy$ plane). We may then reduce the lattice and retain only the neighbour in the $(100)$ direction and scale the overlap $t_{ss\sigma}$ by the appropriate weight (i.e. 2 in this case) as discussed earlier. For the $sp_x$, however, the overlaps in the $(010), (0\bar{1}0), (001)$ and $(00\bar{1})$ are zero. This can be deduced immediately by noticing that the overlap products in these directions are positive for $x > 0$ and negative for $x < 0$. On reduction, although they are related to one another by the symmetry operations of the cubic lattice, we cannot club the neighbours in the $(100)$ and $(\bar{1}00)$ direction with those in the $(010), (0\bar{1}0), (001)$ and $(00\bar{1})$ and will have to put these overlaps to be zero by introducing the symmetry factor $\beta_{LL'}(\ell mn)$.

\[
\langle RL|H|R'L'\rangle \rightarrow \sqrt{\frac{W_{R'}^{R}}{W_{R}}} \, \beta_{LL'}(\ell mn) \, \langle RL|H|R'L'\rangle
\] 

where $\beta_{LL'}(\ell mn)$ is 0 if $R'$ is a symmetric position of $R$ with respect to $L$ and $L'$, otherwise it is 1.
Figure 3 illustrates some of the overlaps involving the $d$ states with $e_g$ and $t_{2g}$ symmetries. Only projections on the $xy$ plane are shown. The positive parts of the lobes are unshaded, while the negative parts of the lobes are shaded.

Figure 3 illustrates some of the overlaps involving the $d$ orbitals in the $e_g$ and $t_{2g}$ symmetries. With this simple reduction procedure the real space part can be reduced to $1/8$-th, keeping only the sites in the $(x > 0, y > 0, z > 0)$ octant and suitably renormalizing the Hamiltonian matrix elements as described earlier.

We still have not exhausted all the symmetries in the full augmented space. As discussed earlier, this space is a direct product of the real (lattice-orbital) space and the configuration space which are disjoint. As a consequence the symmetry operations apply independently to each of them. Since the disorder is homogeneous, the cardinality
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Figure 4. Equivalent configurations on a cubic lattice. Configurations labelled by $\uparrow$ are shown as light spheres, while those labelled by $\downarrow$ are shown as dark spheres. The top twelve configurations are equivalent, while the bottom three are equivalent. The lower right inset shows the numbering scheme used for the lattice sites in the main text.
sequence in configuration space itself has the symmetry of the underlying lattice. To see this, let us look at the figure where we show a part of a cubic lattice where the central site is occupied by a configuration labelled by $\uparrow$; while two of the six nearest neighbours are occupied by configurations labelled by $\downarrow$, and four of them by $\uparrow$s. We note that the twelve configurations in the first three rows of the figure, where the two $\downarrow$s sit at distances $\sqrt{2}$ times the lattice constant, are related to one another by the symmetry operations of the cubic lattice. For example, the second to the fourth configurations on the top row can be obtained from the first by the rotations $\mathcal{R}(\pi/2, \hat{z})$, $\mathcal{R}(\pi, \hat{z})$ and $\mathcal{R}(3\pi/2, \hat{z})$ respectively. The configurations are described by cardinality sequences (as described earlier). The cardinality sequences for the four configurations on the top row of figure are : $\{1\bar{3}\}$, $\{2\bar{3}\}$, $\{\bar{1}\bar{3}\}$ and $\{\bar{2}\bar{3}\}$. From the figure it is easy to see that :

$$\{2\bar{3}\} = \mathcal{R}(\pi/2, \hat{z}) \{1\bar{3}\} ; \quad \{1\bar{3}\} = \mathcal{R}(\pi, \hat{z}) \{1\bar{3}\} \quad \text{and} \quad \{2\bar{3}\} = \mathcal{R}(3\pi/2, \hat{z}) \{1\bar{3}\}$$

This equivalence of the configurations on the lattice is quite independent of the symmetries of the Hamiltonian in real space discussed earlier. Thus, in augmented space, equivalent states are $|R\otimes\{C\}\rangle$ and the set $|\mathcal{R}R\otimes\mathcal{R}\{C\}\rangle$ for all different symmetry operators $\mathcal{R}$ of the underlying lattice. Again the symmetry of the orbitals also rules out the operation of the Hamiltonian at certain symmetric positions discussed earlier.

Once we have defined the reduced Hamiltonian, recursion on the reduced augmented space with starting state $|u_1\rangle = |RL \otimes \{\emptyset\}\rangle$ gives the configuration averaged Green function directly.

In order to give the readers a flavour of the reduction in storage and CPU time, we have carried out two sets of recursion calculations : First, a standard recursion on ordered Ag on a fcc lattice, and second, an augmented space recursion with 15 seed points on a fcc alloy $Ag_{50}Pd_{50}$ both with and without symmetry reduction. The calculations were done on a P4 machine with 1.13 GHz clock speed and 256 MB RAM.
The first point to note is that convergence of the recursion technique is measured by the convergence of the energy moments of the corresponding density of states obtained from recursion \[23, 24\]. These energy moments are defined by:

\[M_n = \int_{-\infty}^{E_F} dE \, E^n n(E)\]

where \(n_e\) is the number of valence electrons. The convergence of these moments with the number of recursion steps \(N\), from which the density of states \(n(E)\) and the Fermi energy \(E_F\) is calculated, is reflected in the errors

\[\Delta M_N^n = |M_N^n - M_N^{n-1}|\]

The figure\[6\] shows the convergence of the first three moments for a Ag\(_{50}\)Pd\(_{50}\) disordered alloy on a fcc lattice in a spd TB-LMTO minimal basis set. It is clear from the figure that for a convergence within a specified error window one has to carry out recursion for a specific large number of steps determined by the error window. Symmetry reduction allows us to carry out recursion exactly over much larger number of steps than is possible for ordinary recursion, given our computational resources. There was a similar

| System | Shells | Steps | Full lattice | Reduced lattice |
|--------|--------|-------|--------------|-----------------|
|        |        |       | Sites | CPU time | Sites | CPU time |
| Model  | 20     | 20    | 24739 | 2.36     | 3385  | 0.54     |
| Ag     | 20     | 30    | 24739 | 47.15    | 3385  | 15.83    |

Table 2. Comparison between system size and CPU time (in seconds) taken for recursion on a P4/256 machine for a full fcc lattice and the reduced lattice in real space.
Symmetry reduction in the TB-LMTO-ASR for random binary alloys

Figure 7. Showing how the number of lattice sites increase with increasing the number of shells in an augmented space map on a disordered binary alloy on a fcc lattice.

observation made in an earlier estimate of the errors of the recursion method [24]. In order to carry out this many recursion steps we have to generate at least this many shells exactly around the starting site in order to avoid surface-like effects.

For the case of the ordered fcc lattice, the figure 6 shows how the size of the map increases as we increase the number of nearest neighbour shells from a starting site, both with and without reduction. Table 2 shows the details of the CPU time and storage space reduction for recursion after applying the symmetry reduction. We have carried out calculations both on a simple s-state system on a fcc lattice, as well as for Ag (with spd minimal TB-LMTO basis) also on a fcc lattice.

For the calculation of a disordered binary alloy on a fcc lattice, the figure 7 shows the enormous decrease in the size of the augmented space map after application of symmetry reduction for a seven nearest neighbour map on a fcc lattice. Table 3 tabulates the reduction in storage and CPU time for a 7 shell, 11 step recursion in augmented space carried over 15 seed points using TB-LMTO potential parameters and structure matrix to built the Hamiltonian. The power of symmetry reduction of storage space is more evident in this example. Such reduction will allow us to stretch our nearest neighbour map up to 9-10 shells, stepping up our accuracy. The figure 8 indicates the

| System       | Shells | Steps | Full lattice | Reduced lattice |
|--------------|--------|-------|--------------|-----------------|
| Model (50-50)| 7      | 7     | 49476        | 6856            |
| Ag_{50}Pd_{50}| 7      | 11    | 49476        | 6856            |

Table 3. Comparison between system size and time (in seconds) taken for recursion (using 15 seed points for Ag_{50}Pd_{50}) on a P4/256 machine for a full fcc lattice and the reduced lattice in augmented space.
Symmetry reduction in the TB-LMTO-ASR for random binary alloys

Figure 8. CPU time taken per recursion step for a $Ag_{50}Pd_{50}$ disordered binary alloy on a fcc lattice in a spd TB-LMTO minimal basis set. The times with and without symmetry reduction are shown.

We have shown that recursion calculations can be carried out much faster and for many more recursive steps exactly, if we perform the recursion on a subspace of the original augmented space reduced by using the symmetries of both the underlying lattice and random configurations on the lattice. This will allow us to obtain results for disordered binary alloys with enhanced accuracy required for first-principles, self-consistent, density functional based calculations. In this communication we have described the details of the implementation of this symmetry reduction and the modifications required in the standard recursion method. We propose using the symmetry reduced version of the augmented space recursion in our future work on disordered alloys.

4. Conclusions

We have shown that recursion calculations can be carried out much faster and for many more recursive steps exactly, if we perform the recursion on a subspace of the original augmented space reduced by using the symmetries of both the underlying lattice and random configurations on the lattice. This will allow us to obtain results for disordered binary alloys with enhanced accuracy required for first-principles, self-consistent, density functional based calculations. In this communication we have described the details of the implementation of this symmetry reduction and the modifications required in the standard recursion method. We propose using the symmetry reduced version of the augmented space recursion in our future work on disordered alloys.

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