Operads as polynomial 2-monads

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Abstract. In this article we give a construction of a polynomial 2-monad from an operad and describe the algebras of the 2-monads which then arise. This construction is different from the standard construction of a monad from an operad in that the algebras of our associated 2-monad are the categorified algebras of the original operad. Moreover it enables us to characterise operads as categorical polynomial monads in a canonical way. This point of view reveals categorical polynomial monads as a unifying environment for operads, Cat-operads and clubs. We recover the standard construction of a monad from an operad in a 2-categorical way from our associated 2-monad as a coidentifier of 2-monads, and understand the algebras of both as weak morphisms of operads into a Cat-operad of categories. The relevance of this to the theory of internal algebras is discussed.

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1. Introduction

In contemporary mathematics there has been a proliferation of operadic notions. These include cyclic operads, modular operads, dioperads, properads and so on, with the basic combinatorics underpinning these notions being more involved than that of the standard operads that arose originally in algebraic topology in the 1970’s. One of the many contributions of Batanin and Berger in [1], is to exhibit these contemporary operadic notions as algebras of very particular standard operads. In this article we use unadorned name “operad” for what are commonly

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referred to as “coloured symmetric operads of sets” or also as “symmetric multicategories”. More precisely, Batanin and Berger exhibited many of these contemporary operadic notions as algebras of Σ-free operads, an operad being Σ-free when its symmetric group actions admit no fixed points.

An operad \( T \) with set of colours \( I \) determines a monad we denote as \( T/\Sigma \) on the category \( \text{Set}/I \) whose algebras are the algebras of the operad \( T \). When \( T \) is Σ-free this monad is a polynomial monad in the sense of [5]. In fact as explained by Kock [15] and Szawiel-Zawadowski [21], finitary polynomial monads may be identified with Σ-free operads. In general, polynomial monads are examples of cartesian monads, so may be applied to internal categories, and in this way one may regard \( T/\Sigma \) as a 2-monad on \( \text{Cat}/I \). Thus any contemporary operadic notion determines a 2-monad, and so the rich theory of 2-dimensional monad theory [3] becomes applicable in these contexts. This is part of the technology which underpins [1], and which is developed further in this paper and [2, 23].

In this article we give an alternative construction of a polynomial 2-monad from an operad and describe the algebras of the 2-monads which then arise. Our construction does not require Σ-freeness, and in the case of a Σ-free operad \( T \) with set of colours \( I \), produces a 2-monad on \( \text{Cat}/I \) which is different from \( T/\Sigma \). In the general case we give an alternative 2-categorical construction of \( T/\Sigma \) from \( T \), and then establish that this construction restricts to the world of polynomial 2-monads in the Σ-free case.

This alternative construction comes from a characterisation of operads and related notions in terms of polynomials [5, 26] in \( \text{Cat} \). Recall that a polynomial monad in \( \text{Cat} \) is a monad in a certain bicategory \( \text{PolyCat} \) whose objects are categories, the underlying endomorphism \( I \to I \) of which is a diagram as on the left

\[
\begin{array}{ccc}
I & \xleftarrow{s} & E & \xrightarrow{p} & B & \xrightarrow{t} & I \\
1 & \xleftarrow{\mathbb{P}} & \mathbb{P} & \xrightarrow{U^p} & \mathbb{P} & \xrightarrow{t} & 1
\end{array}
\]

in which \( p \) is an exponentiable functor. In particular one has the polynomial monad \( \text{sm} \) indicated on the right in which \( \mathbb{P} \) is a skeleton of the category of finite sets and bijections, and a morphism from the former to the latter consists of the functors \( e \) and \( b \) fitting into a commutative diagram

\[
\begin{array}{ccc}
I & \xleftarrow{s} & E & \xrightarrow{p} & B & \xrightarrow{t} & I \\
1 & \xleftarrow{\mathbb{P}} & \mathbb{P} & \xrightarrow{U^p} & \mathbb{P} & \xrightarrow{t} & 1 \\
& & e & \xrightarrow{pb} & b
\end{array}
\]

and are compatible with the monad structures. By theorem 3.3, remarks 3.4 and 3.5, and proposition 6.3 one may identify

- Operads as situations (1) in which \( I \) is discrete and \( b \) is a discrete fibration.
- Σ-free operads as situations (1) in which \( I \) is discrete, \( b \) is a discrete fibration and \( B \) is equivalent to a discrete category.
- \( \text{Cat} \)-operads as situations (1) in which \( I \) is discrete and \( b \) has the structure of a cloven split fibration.
- Operadic categories in the sense of Toën [22] as situations (1) in which \( I \) is discrete and \( b \) has the structure of a cloven fibration.
- Clubs in the sense of Kelly [10, 11] as situations (1) in which \( I = 1 \).
Our construction of a 2-monad from an operad regards an operad $T$ in this way in which $I$ is the set of colours, and then the usual construction [5, 26] of a polynomial functor from a polynomial gives a 2-monad on $\text{Cat}/I$. The basic theory of polynomials and polynomial functors from [5, 26] is recalled in section 2.

The 2-monad on $\text{Cat}/I$ just described is also denoted as $T$. We find this convention to be most convenient, but when using it one must be aware that the conventional $\text{Cat}$-valued algebras of the operad $T$ are the strict algebras of the 2-monad $T/\Sigma$, whereas the algebras of the 2-monad $T$ correspond to categorified algebras of the operad $T$. For example when $T$ is the terminal operad $\text{Com}$, a strict algebra for the 2-monad $T$ is a symmetric strict monoidal category, whereas a strict algebra for $T/\Sigma$ is a commutative monoid in $\text{Cat}$.

Like any 2-monad, $T$ has different types of algebra (lax, pseudo and strict), different types of algebra morphism (lax, colax, pseudo and strict) and thus a variety of different 2-categories of algebras depending on which types of algebras and algebra morphisms one is interested in. By theorem 4.8 $T$-algebras admit an explicit description as the appropriately weak morphisms of operads $T \to \text{Cat}$, where $\text{Cat}$ is a canonical $\text{Cat}$-enriched operad whose objects are categories. That is, the lax, pseudo and strict algebras of the 2-monad $T$ are lax, pseudo and strict morphisms of operads $T \to \text{Cat}$ in the sense of definitions 4.1 and 4.2. In the case where $T$ is a category, that is when all its operations are of arity 1, this description of the algebras of $T$ is well-known and goes back at least to [3]. In this case the 2-monad $T$ is the 2-monad on $\text{Cat}/I$ whose 2-category of strict algebras and strict morphisms is the functor 2-category $[T, \text{Cat}]$, and a lax or pseudo algebra is exactly a lax or pseudo functor $T \to \text{Cat}$.

Similarly one has characterisations of the various types of $T$-algebra morphism in theorem 4.12 and $T$-algebra 2-cells in theorem 4.15. In particular for any operad $T$ and any symmetric monoidal category $\mathcal{V}$, algebras of $T$ in $\mathcal{V}$ can be seen as lax $T$-algebra morphisms in a canonical way. As explained at the end of section 4 the via the central examples of [1], one exhibits the categories cyclic operads, modular operads, dioperads, properads and so on, in a symmetric monoidal category, in this way.

To understand the relationship between the 2-monads $T$ and $T/\Sigma$ on $\text{Cat}/I$ for a given operad $T$ with set of colours $I$, one begins by thinking about the algebras. One feature of the notion of lax, pseudo or strict morphism $H : T \to \text{Cat}$ alluded to above is that $H$ is not equivariant in the strictest sense, but rather that it is so up to coherent isomorphisms which are called the symmetries of $H$. When these symmetries are identities, the lax morphism is said to be commutative, and it is the commutative strict morphisms $T \to \text{Cat}$ which correspond to the algebras of $T/\Sigma$. In other words strict $T/\Sigma$-algebras are included amongst strict $T$-algebras, and this inclusion is exactly the inclusion of the commutative strict morphisms $T \to \text{Cat}$ amongst the general strict morphisms.

The standard construction of $T/\Sigma$ is via a formula which involves quotienting out by the symmetric group actions of $T$. In this article this quotienting is carried in a 2-categorical way. Starting from the 2-monad $T$ definition 5.6 provides a 2-cell of 2-monads $\alpha_T$ as in

$$T^{[1]}_\Sigma \xrightarrow{d_T} T \xrightarrow{q_T} T/\Sigma$$
and then $q_T$ is the universal morphism of 2-monads which post-composes with $\alpha_T$ to give an identity. In the language of 2-category theory, $q_T$ is the reflexive coidentifier of $\alpha_T$ in $\text{Mnd}(\text{Cat}/I)$. The algebras of $T/\Sigma$ defined in this way are identified with commutative operad morphisms $T \to \text{Cat}$ in theorem 5.16, and so our construction of $T/\Sigma$ coincides with the standard one. Moreover as explained in section 6 when $T$ is a $\Sigma$-free operad one can witness this quotienting process as taking place completely in the world of polynomials, and this is why $T/\Sigma$ is a polynomial monad when $T$ is $\Sigma$-free.

Notation, terminology and background. We assume a basic familiarity with some of the elementary notions of 2-category theory – basic 2-categorical limits such as cotensors, comma objects and isocomma objects; the calculus of mates as explained in [14]; and the basic notions 2-dimensional monad theory which one can find for instance in [3, 16]. This article is a sequel to [26], and so one can find an exposition of many background notions relevant here, such as fibrations and their definition internal to any 2-category, in [26] in a way that is notationally and terminologically compatible with this article. However some effort has been made to recall important background as needed for the convenience of the reader. For instance one finds the definition of the various types of algebra of a 2-monad recalled in section 4 just before the details of these definitions are needed in this work. As in [26] we denote by $[n]$ the ordinal $\{0 < \ldots < n\}$ regarded as a category.

We denote by $\text{Cat}$ the 2-category of small categories, and sometimes make use of a 2-category $\text{CAT}$ of large categories, which include standard categories of interest, like $\text{Set}$ as objects. Such size issues can and have been axiomatised [20, 25], but for the most part in this paper we pay little attention to such matters beyond the notation just described.

2. Polynomial functors

Composition with an arrow $f : X \to Y$ in a category $\mathcal{E}$ defines the effect on objects of the functor $\Sigma_f : \mathcal{E}/X \to \mathcal{E}/Y$. When $\mathcal{E}$ has pullbacks $\Sigma_f$ has a right adjoint denoted $\Delta_f$, given on objects by pulling back along $f$. When $\Delta_f$ has a right adjoint, $f$ is said to be exponentiable and this further right adjoint is denoted $\Pi_f$. The exponentiable arrows in a category with pullbacks are closed under composition, and stable by pullback along arbitrary maps. A category which has all pullbacks, and all of whose arrows is exponentiable is said to be locally cartesian closed. While there are many examples of such categories, for instance any elementary topos is locally cartesian closed, our main example $\mathcal{E} = \text{Cat}$ is not. However, in this case one has the combinatorial characterisation of exponentiable functors as Giraud-Conduché fibrations [4, 6, 18]. In particular Grothendieck fibrations and Grothendieck opfibrations are Giraud-Conduché fibrations.

In elementary terms, the effect of $\Pi_f$ on objects is to take distributivity pullbacks along $f$ in the sense to be recalled now from [26]. Given $g : W \to X$, a
pullback around \((f,g)\) consists of \((p,q,r)\) as on the left in

\[
\begin{array}{ccc}
P & \xrightarrow{q} & Q \\
\downarrow{p} & & \downarrow{g} \\
W & \xrightarrow{\text{pb}} & r \\
\downarrow{f} & & \downarrow{\text{pb}} \\
X & \xrightarrow{\text{pb}} & Y \\
\end{array}
\]

\[
\begin{array}{ccc}
P & \xrightarrow{q} & Q \\
\downarrow{p} & & \downarrow{\text{pb}} \\
W & \xrightarrow{\Delta_f \Pi_f(g)} & r \\
\downarrow{g} & & \downarrow{\Pi_f(g)} \\
X & \xrightarrow{f} & Y \\
\end{array}
\]

such that the morphisms \((gp,f,r,q)\) form a pullback square as indicated. A morphism \((p,q,r)\) consists of morphisms \((p',q',r')\) such that \(p = p'\alpha, \beta q = q'\alpha\) and \(r = r'\beta\). A distributivity pullback around \((f,g)\) is a terminal object in the category of pullbacks around \((f,g)\) just described. A general such is denoted as in the middle of the previous display. The connection with \(\Pi_f\) is indicated on the right in the previous display, in which \(\varepsilon^f\) is the counit of \(\Delta_f \dashv \Pi_f\) and \(\varepsilon^f_g\) is its component at \(g\). Explicitly, the universal property of a distributivity pullback says that given \((p',q',r')\) as in

\[
\begin{array}{ccc}
P & \xrightarrow{q} & Q \\
\downarrow{p} & & \downarrow{g} \\
W & \xrightarrow{\text{pb}} & r \\
\downarrow{f} & & \downarrow{\text{pb}} \\
X & \xrightarrow{\Delta_f \Pi_f(g)} & Y \\
\end{array}
\]

making the square with boundary \((gp,f,r',q')\) a pullback, there exist \(\alpha\) and \(\beta\) as shown unique such that \(p\alpha = p', q\alpha = \beta q'\) and \(r\beta = r'\).

In practice one is often interested in obtaining an explicit description of \(Q\) and \(r\) in terms of the generating data \((f,g)\) of a distributivity pullback. For instance when the ambient category \(E\) is \(\text{Set}\) then an element \(z\) of \(Q\) over \(y \in Y\) as on the left

\[
\begin{array}{ccc}
[0] & \xrightarrow{z} & Q \\
\downarrow{y} & & \downarrow{r} \\
Y & \xrightarrow{f^{-1}\{y\}} & W \\
\end{array}
\]

amounts, by the adjunction \(\Delta_f \dashv \Pi_f\), to the morphism \(h\) on the right in the previous display where \(i_y\) is the inclusion. Thus one can identify elements of \(Q\) as pairs \((y,h)\), and then \(r\) is given by \(r(y,h) = y\).

We now describe a similar explicit description of distributivity pullbacks in \(\text{Cat}\) in the case where \(f\) is a discrete opfibration. To this end recall that for a discrete opfibration \(f : X \to Y\) one has the corresponding functor \(\tilde{f} : Y \to \text{Set}\) whose effect on objects is \(y \mapsto f^{-1}\{y\}\), and whose lax colimit is \(X\). The data of the lax colimit cocone consists of the inclusions of fibres \(i_y : f^{-1}\{y\} \to X\) for all \(x \in X\), and lax naturality 2-cells \(i_\alpha : i_{y_1} \to i_{y_2} \tilde{f}(\alpha)\) for all \(\alpha : y_1 \to y_2\) in \(Y\). Another way to organise this information uses the Fam construction. For a category \(C\), the category \(\text{Fam}(C)\) has as objects pairs \((I,h)\) where \(I\) is a set regarded as a discrete category, and \(h : I \to C\) is a functor. A morphism \((I,h) \to (J,k)\) consists of \((f,\phi)\)
where $f : I \to J$ is a function and $\phi : h \to kf$ is a natural transformation. Then the fibres of $f$ and the above lax colimit cocone organise to form a functor\(^1\)

\[ \mathcal{F} : Y \to \text{Fam}(X) \quad y \mapsto (f^{-1}\{y\}, i_y) \quad \alpha : y_1 \to y_2 \mapsto (\tilde{f}(\alpha), i_\alpha). \]

**Lemma 2.1.** In a distributivity pullback as on the left

\[
\begin{array}{ccc}
P & \xrightarrow{g} & W \xrightarrow{f} X \\
| & \downarrow & | \downarrow \\
Q & \xrightarrow{r} & Y
\end{array}
\quad \begin{array}{ccc}
Q & \xrightarrow{pb} & \text{Fam}(W) \\
| & \downarrow & | \downarrow \\
Y & \xrightarrow{\mathcal{F}} & \text{Fam}(X)
\end{array}
\]

in \textbf{Cat} in which $f$ is a discrete opfibration, $r$ and $R$ can be described explicitly by the pullback on the right.

**Proof.** Since discrete opfibrations are exponentiable functors the distributivity pullback exists, and one just needs to use the adjunction $\Delta_f \dashv \Pi_f$ to unpack the explicit description and match it up with the pullback on the right. On objects the situation is as with the case $\mathcal{E} = \text{Set}$, and so an object of $Q$ may be identified with $(y, h)$ where $y \in Y$ and $h : f^{-1}\{y\} \to W$ such that $gh = i_y$, which is exactly an object of the pullback on the right. An arrow of $Q$ amounts to a functor $\alpha : [1] \to Q$, and thus a choice of arrow $\alpha : y_1 \to y_2$ in $Y$ codified itself as a functor $\alpha : [1] \to Y$, together with $\beta : [1] \to Q$ such that $r\beta = \alpha$. Pulling back $\alpha$ along $f$ gives a category whose object set is the disjoint union $f^{-1}\{y_1\} \coprod f^{-1}\{y_2\}$, with one non-identity morphism for each element of $x \in f^{-1}\{y_1\}$. The domain of the morphism corresponding to $x$ is $x$ itself, and its codomain is $\tilde{f}(\alpha)(x)$. Thus by the adjunction $\Delta_f \dashv \Pi_f$, a morphism of $Q$ amounts to a morphism of the pullback. \(\square\)

The bicategory $\text{Poly}_\mathcal{E}$ of polynomials in $\mathcal{E}$ was first described explicitly in the locally cartesian closed case in [5], and for a general category with pullbacks in [26]. An object of $\text{Poly}_\mathcal{E}$ is an object of $\mathcal{E}$. An arrow $I \to J$ in $\text{Poly}_\mathcal{E}$ is a polynomial in $\mathcal{E}$ from $I$ to $J$, which by definition is a diagram as on the left

\[
\begin{array}{ccc}
I & \xleftarrow{s} & E & \xrightarrow{p} & B & \xrightarrow{t} & J \\
| & | & | & | & | & | \\
E_1 & \xrightarrow{p_1} & B_1 & \xrightarrow{t_1} & J \\
E_2 & \xrightarrow{p_2} & B_2 & \xrightarrow{t_2} & J
\end{array}
\]

in $\mathcal{E}$ in which the middle map $p$ is exponentiable. A 2-cell $f : (s_1, p_1, t_1) \to (s_2, p_2, t_2)$ in $\text{Poly}_\mathcal{E}$ is a diagram as on the right in the previous display, and we call the morphisms $f_1$ and $f_2$ the *components* of $f$. In elementary terms the process of forming the horizontal composite $(s_3, p_3, t_3) = (s_2, p_2, t_2) \circ (s_1, p_1, t_1)$ of polynomials

---

\(^1\)As explained in section 5 of [24], $\mathcal{F}$ is the Fam-generic factorisation of $\tilde{f} : Y \to \text{Set} = \text{Fam}(1)$. 
is encapsulated by the commutative diagram

\[
\begin{array}{c}
E_3 \xymatrix{ \ar[r]^{p_1} & F } \\
\ar[d]^{s_3} & B_1 \times_J E_2 \ar[r]^{dp_b} & B_2 \ar[d]^{t_3} \\
I \ar[r]_{s_1} & E_1 \ar[r]_{p_1} & B_1 \ar[r]_{t_1} & J \ar[r]_{s_2} & E_2 \ar[r]_{p_2} & B_2 \ar[r]_{t_2} & K.
\end{array}
\]

At this level of generality there is, by theorem 3.2.6 of [26], a homomorphism

\[ P_\mathcal{E} : \text{Poly}_\mathcal{E} \rightarrow \text{CAT} \quad I \mapsto \mathcal{E}/I \]

of bicategories with object map as indicated. The effect of \( P_\mathcal{E} \) on arrows is to send the polynomial \((s,p,t)\) to the composite functor \( \Sigma t \Pi p \Delta s : \mathcal{E}/I \rightarrow \mathcal{E}/J \). In elementary terms \( P_\mathcal{E}(s,p,t) \) has object map as indicated on the left in

\[
\begin{array}{c}
X_3 \xymatrix{ \ar[r] & X_4 } \\
\ar[d] & \ar[d] \ar[d] \\
X \xymatrix{ \ar[r]_{dp_b} & X_2 } \\
\ar[d]^{s} & \ar[d]^{p} & \ar[d]^{t} \\
I & E \ar[r]^{p} & B \ar[r]^{t} & J
\end{array}
\]

\[
\begin{array}{c}
X_3 \xymatrix{ \ar[r]^{dp_b} & X_4 } \\
\ar[d]^{p_b} & \ar[d]^{p_b} \ar[d] \ar[d] \\
X_2 \ar[r]_{dp_b} & X_3 \\
\ar[d]^{s} & \ar[d]^{s} & \ar[d]^{s} \\
I & E \ar[r]_{dp_b} & B \ar[r]_{dp_b} & J
\end{array}
\]

Given a morphism \( f \) in \( \mathcal{E}/I \), the morphism \( X_4 \rightarrow Y_4 \) on the right is the effect of \( P_\mathcal{E}(s,p,t) \) on \( f \) in elementary terms. To obtain it one first induces \( X_2 \rightarrow Y_2 \) using the pullback that defines \( Y_2 \), and then both the maps \( X_3 \rightarrow Y_3 \) and \( X_4 \rightarrow Y_4 \) are induced by the universal property of the distributivity pullback \((Y_3,Y_4,B,E,Y_2)\). The elementary description of the 2-cell map of \( P_\mathcal{E} \) is given by

\[
\begin{array}{c}
X_2 \xymatrix{ \ar[r]^{dp_b} & X_3 } \\
\ar[d]^{p_b} & \ar[d]^{p_b} \ar[d] \ar[d] \\
X \ar[r]_{dp_b} & E \ar[r]_{dp_b} & B \ar[r]_{dp_b} & J
\end{array}
\]

in which the component at \( x \) of \( P_\mathcal{E}(f_1,f_2) \) is the induced map \( X_4 \rightarrow X'_4 \).

**Example 2.2.** As an illustration consider the case of a polynomial \((s,p,t)\) as above in \( \text{Cat} \), in which \( p : E \rightarrow B \) is a discrete opfibration, and write \( T : \text{Cat}/I \rightarrow \text{Cat}/I' \).
\[ \textbf{Cat}/J \] for the polynomial functor \( P_{\text{Cat}}(s, p, t) \). Then for \( X \to I \) in \( \textbf{Cat}/I \), \( TX \) is formed as on the left

\[
\begin{align*}
X & \xleftarrow{p} E \xrightarrow{\phi} B \xrightarrow{p} \text{Fam}(X \times_I E) \xrightarrow{\text{Fam}(I)} \text{Fam}(E) \\
I & \xleftarrow{s} I \xrightarrow{p} E \xrightarrow{\phi} B \xrightarrow{p} \text{Fam}(X) \xrightarrow{\text{Fam}(I)} \text{Fam}(I)
\end{align*}
\]

and so in view of the fact that the Fam construction preserves pullbacks and lemma 2.1, one has pullbacks in \( \textbf{CAT} \) as on the right. Unpacking the composite pullback on the right, one finds that \( TX \) has the following explicit description. An object is a pair \((b, h)\) where \( b \) is an object of \( B \) and \( h : p^{-1}\{b\} \to X \) whose composite with \( X \to I \) is \( s_i b \), where \( i_b : p^{-1}\{b\} \to E \) is the inclusion of the fibre. A morphism \((b_1, h_1) \to (b_2, h_2)\) is a pair \((\beta, \gamma)\) where \( \beta : b_1 \to b_2 \) is in \( B \) and \( \gamma \) is a natural transformation as on the left

\[
\begin{align*}
\xymatrix{ p^{-1}\{b_1\} \ar[r]^-{\bar{\beta}} & p^{-1}\{b_2\} \ar[d]^-{h_1} & p^{-1}\{b_1\} \ar[r]^-{\bar{\beta}} & p^{-1}\{b_2\} \ar[d]^-{h_2} \\
X \ar[u]^-{b_1} \ar[r]_-{h_2} & E \ar[u]^-{s} & E \ar[u]^-{s} \ar[r]_-{i_{b_2}} & E \ar[u]^-{s} \ar[r]_-{i_{b_1}} & E \ar[u]^-{s}
\end{align*}
\]

satisfying the equation, in which \( \bar{\beta}(\beta) \) and \( i_{\beta} \) are as described prior to lemma 2.1.

In the cases of interest in this article, \( I \) is typically discrete, in which case this last equation is satisfied automatically.

A span in \( \mathcal{E} \) as on the left

\[
\begin{align*}
I & \xleftarrow{s} B \xrightarrow{t} J \\
I & \xleftarrow{s} B \xrightarrow{1_B} B \xrightarrow{t} J
\end{align*}
\]

is identified with a polynomial as on the right, and the composition of polynomials generalises the usual pullback composition of spans. In particular \( f : I \to J \) in \( \mathcal{E} \) determines the spans \( f^* \) and \( f_* \)

\[
\begin{align*}
I & \xleftarrow{1_I} I \xrightarrow{f} J \\
J & \xrightarrow{f} I \xrightarrow{1_I} I \\
J & \xrightarrow{f} I \xrightarrow{1_I} I
\end{align*}
\]

as on the left and middle respectively, and one has an adjunction \( f^* \dashv f_* \) in \( \text{Span}_E \) and thus also in \( \text{Poly}_E \). The counit \( c_f \) of this adjunction in \( \text{Poly}_E \) is given by the diagram on the right in the previous display.

For a category \( \mathcal{E} \) with pullbacks one also has, as described in [5] in the locally cartesian closed case, the categories \( \text{PolyEnd}_E \) and \( \text{PolyMnd}_E \) of polynomial endomorphisms and polynomial monads respectively. An object of \( \text{PolyEnd}_E \) is
a pair \((I, P)\) where \(I\) is an object of \(E\) and \(P : I \to I\) is in \(\text{Poly}_E\). A morphism \((I, P) \to (J, Q)\) is a pair \((f, \phi)\) where \(f : I \to J\) is in \(E\) and \(\phi : f^* \circ P \circ f_* \to Q\) is in \(\text{Poly}_E\). Given morphisms \((f, \phi) : (I, P) \to (J, Q)\) and \((g, \gamma) : (J, Q) \to (K, R)\), the underlying arrow of \(E\) of the composite \((g, \gamma)(f, \phi)\) is \(gf\), and the 2-cell datum of this composite is given by

\[
g^* \circ f^* \circ P \circ f_* \circ g_* \xrightarrow{g^* \circ \phi \circ g^*} g^* \circ Q \circ g_* \xrightarrow{\gamma} R.
\]

In more elementary terms, writing \(P = (s, p, t)\) and \(Q = (s', p', t')\), a morphism \((f, \phi) : (I, P) \to (J, Q)\) of \(\text{PolyEnd}_E\) amounts to a commutative diagram

\[
\begin{array}{ccc}
I & \xrightarrow{s} & E \\
\downarrow{f} & & \downarrow{p} \\
J & \xrightarrow{s'} & E'
\end{array}
\quad
\begin{array}{ccc}
P & \xrightarrow{b} & B \\
\downarrow{f} & & \downarrow{t} \\
Q & \xrightarrow{b'} & B'
\end{array}
\]

(2)

because \(f^* \circ P \circ f_* = (fs, p, ft)\), and the composition just described amounts to stacking such diagrams vertically. The various mates of \(\phi : f^* \circ P \circ f_* \to Q\) with respect to \(f^* \circ f_*\) are denoted

\[
\begin{align*}
\phi^e : f^* \circ P & \to Q \\
\phi^f : P & \to f^* \circ Q \\
\phi^i : P & \to f_* \circ Q \circ f^*.
\end{align*}
\]

Note in particular that when \(Q\) underlies a monad on \(I\) in \(\text{Poly}_E\), then the composite \(f_* \circ Q \circ f^*\) underlies a monad on \(I\).

An object of \(\text{PolyMnd}_E\) is again a pair \((I, P)\) with \(P\) this time a monad on \(I\) in \(\text{Poly}_E\), and we shall often adopt the abuse of referring to both the endomorphism and the monad as \(P\). A morphism \((I, P) \to (J, Q)\) in \(\text{PolyMnd}_E\) is a morphism \((f, \phi)\) in \(\text{PolyEnd}_E\), together with the condition that \(\phi : P \to f_* \circ Q \circ f^*\) is a morphism of monads on \(I\). This last condition admits reformulations in the language of [19], namely that \((f^*, \phi^e)\) is a monad opfunctor, or equivalently, that \((f_*^*, \phi^i)\) is a monad functor.

The situation of interest for us in this article is when \(E = \text{Cat}\), and in this case polynomial monads give rise to 2-monads on slices of \(\text{Cat}\), and the natural transformations \(\text{P}_{\text{Cat}}(\phi^e)\) and \(\text{P}_{\text{Cat}}(\phi^i)\) considered above are 2-natural. This is so because one has a 2-categorical version of the theory of polynomials and polynomial functors described in section 4 of [26], and applied to the 2-category \(\text{Cat}\) enables one to see \(\text{P}_{\text{Cat}}\) as taking values in \(\text{2-CAT}\). In fact more is true, because both \(\text{P}_{\text{Cat}}\) and \(\text{2-CAT}\) form degenerate sorts of tricategories called 2-bicategories, which are just like bicategories except that the homs are strict 2-categories instead of mere categories. Moreover \(\text{P}_{\text{Cat}}\) is a homomorphism of 2-bicategories. In elementary terms, a 3-cell \((f_1, f_2) \to (g_1, g_2)\) in \(\text{Poly}_{\text{Cat}}\) consists of \((\alpha, \beta)\) as in

making the diagram commute, that is, such that \(s_2 \beta = \text{id}\), \(\beta p_2 = \alpha p_1\) and \(t_2 \alpha = \text{id}\).

The basic example of a 2-monad on \(\text{Cat}\) arising from a polynomial in \(\text{Cat}\) is the 2-monad for symmetric monoidal categories, and was described in detail in
section (5.2) of [26]. This polynomial is denoted as $\text{sm}$, and its underlying endomorphism in $\text{Poly}_{\text{Cat}}$ is

$$
\begin{array}{ccc}
1 & \xrightarrow{\alpha} & P \\
\downarrow^{p} & & \downarrow^{P} \\
P & \xrightarrow{T} & P
\end{array}
$$

where $P$ is the category of natural numbers and permutations (that is, a skeleton of the category of finite sets and bijections), $P_*$ is the corresponding skeleton of finite pointed sets and base point preserving bijections, and $U^P$ is the forgetful functor. The resulting 2-monad $P_{\text{Cat}}(\text{sm})$ on $\text{Cat}$ is denoted as $\text{Sm}$. As explained in [26], the properties on $U^P$ ensure that $\text{Sm}$ has good formal properties – it is familial, opfamilial and sifted colimit preserving. Taking into account also the canonical linear order on each fibre of $U^P$ leads, as in [26] proposition (5.2.1), to a conceptual explanation of $\text{sm}$'s monad structure.

3. Operads as polynomial monads

We begin by recalling some basic definitions and fixing our notation and terminology. It will often be convenient to denote a typical element $(x_1, \ldots, x_n)$ of a cartesian product $\prod_{i=1}^{n} X_i$ of sets as $(x_i)_{1 \leq i \leq n}$, or as $(x_i)_1$ when $n$ is understood or when we wish it to be implicit. Moreover we denote by $\Sigma_n$ the group of permutations of $\{1, \ldots, n\}$.

A collection $T$ consists of a set $I$ whose elements are called the objects or colours of $X$, and for each pair $((i_j)_{1 \leq j \leq n}, i)$ consisting of a sequence $(i_j)_j$ of elements of $I$ and a single element $i \in I$, one has a set $T((i_j)_j;i)$ whose elements are called arrows of $T$ with source $(i_j)_j$ and target $i$, and a typical element may be denoted as $\alpha : (i_j)_j \rightarrow i$. Furthermore given an arrow $\alpha : (i_j)_{1 \leq j \leq n} \rightarrow i$ and a permutation $\rho \in \Sigma_n$, one has an arrow $\alpha \rho : (i_{\rho(j)})_j \rightarrow i$, this assignation being functorial in the sense that $\alpha 1_n = \alpha$ and $(\alpha \rho_1) \rho_2 = \alpha(\rho_1 \rho_2)$ for all $\rho_1, \rho_2 \in \Sigma_n$.

To any collection $T$ whose set of colours is $I$ we associate a morphism

$$
\begin{array}{ccc}
I & \xrightarrow{s_T} & E_T \\
\downarrow & \downarrow & \downarrow \\
1 & \xrightarrow{\beta_T} & I
\end{array}
$$

and

$$
\begin{array}{ccc}
P_* & \xrightarrow{p_T} & B_T \\
\downarrow & \downarrow & \downarrow \\
P & \xrightarrow{t_T} & 1
\end{array}
$$

(3)

of $\text{Poly}_{\text{Endcat}}$ as follows. Denoting by $T_n$ the set of arrows of $T$ whose source is a sequence of length $n$, $n \mapsto T_n$ is the effect on objects of a functor $P^{op} \rightarrow \text{Set}$, and the corresponding discrete fibration is $b_T : B_T \rightarrow P$. In explicit terms an object of $B_T$ is an arrow $\alpha : (i_j)_j \rightarrow i$ of $T$, and an arrow $\alpha \rightarrow \beta$ of $B_T$ is a permutation $\rho$ such that $\alpha = \beta \rho$. An object of $E_T$ is a pair $\alpha : (i_j)_{1 \leq j \leq n} \rightarrow i$ is an arrow of $T$ and $1 \leq j \leq n$, and an arrow $\alpha : (i_j)_{1 \leq j \leq n} \rightarrow i$ is a permutation $\rho$ such that $\alpha = \beta \rho$ and $\rho j = k$. Thus a typical arrow of $E_T$ can be written as $\rho : (\alpha \rho, j) \rightarrow (\alpha, j)$. The object maps of $s_T, p_T, t_T, b_T$ and $e_T$ are

$$
\begin{array}{ccc}
i & \xrightarrow{i_{(n,j)}} & (n,j) \\
\downarrow & \downarrow & \downarrow \\
\beta_T & \xrightarrow{e_T} & n
\end{array}
$$

and

$$
\begin{array}{ccc}
i_{(\alpha,j)} & \xrightarrow{p_T} & \alpha \\
\downarrow & \downarrow & \downarrow \\
\beta_T & \xrightarrow{e_T} & i
\end{array}
$$

and

$$
\begin{array}{ccc}
i & \xrightarrow{i_{(n,j)}} & (n,j) \\
\downarrow & \downarrow & \downarrow \\
\beta_T & \xrightarrow{e_T} & n
\end{array}
$$

and

$$
\begin{array}{ccc}
i_{(\alpha,j)} & \xrightarrow{p_T} & \alpha \\
\downarrow & \downarrow & \downarrow \\
\beta_T & \xrightarrow{e_T} & i
\end{array}
$$

and

$$
\begin{array}{ccc}
i & \xrightarrow{i_{(n,j)}} & (n,j) \\
\downarrow & \downarrow & \downarrow \\
\beta_T & \xrightarrow{e_T} & n
\end{array}
$$
in which \( n \) is the length of the domain sequence of \( \alpha \), the arrow maps are defined analogously, and the pullback square is easily verified. Since \( U^p \) is a discrete fibration with finite fibres, so is \( p_T \) since such properties on a functor are pullback stable. Thus \( p_T \) is an exponentiable functor. We denote by \( P_T \) the polynomial \((s_T, p_T, t_T)\), and by \( NT : P_T \to \text{sm} \) the morphism \((b_T, e_T)\) of \( \text{PolyMnd}_{\text{Cat}} \).

Let \( S \) and \( T \) be collections with object sets \( I \) and \( J \) respectively. Then a morphism \( F : S \to T \) consists of a function \( f : I \to J \) between object sets and for each \( (i_j)_{1 \leq j \leq n, i} \), a function \( F((i_j)_{j,i}) : S((i_j); i) \to T((f(i_j)); f i) \). These arrow mappings must be equivariant in the sense that given \( \alpha : (i_j)_{1 \leq j \leq n} \to i \) in \( S \) and a permutation \( \rho \in \Sigma_n \), \((f \alpha \rho) = f(\alpha \rho)\). The category of collections and their morphisms is denoted \( \text{Coll} \).

The assignment of a morphism \((3)\) to a collection is the object map of a functor

\[
\mathcal{N} : \text{Coll} \to \text{PolyEnd}_{\text{Cat}} \downarrow \text{sm}
\]

whose arrow map we shall now describe. Given a morphism \( F : S \to T \) of collections the functor \( F_1 : B_S \to B_T \) on objects acts as the arrow map of \( F \), and sends \( \rho : \alpha \rho \to \alpha \) to \( \rho : (F \alpha) \rho \to F \alpha \). Clearly one has \( F_1 b_T = b_S \). The functor \( F_2 : E_S \to E_T \) sends \( (\alpha, j) \) and \( \rho : (\alpha \rho, j) \to (\alpha \rho j) \) to \( (F \alpha, j) \) and \( \rho : (F \alpha \rho j) \to (F \alpha \rho j) \) respectively. Clearly one has \( F_2 e_T = e_X \) and that \((f, F_1, F_2)\) are the components of a morphism \((I, (s_S, p_S, t_S)) \to (J, (s_T, p_T, t_T))\) of \( \text{PolyEnd}_{\text{Cat}} \).

**Proposition 3.1.** The functor \( \mathcal{N} \) restricts to an equivalence between \( \text{Coll} \) and the full subcategory of \( \text{PolyEnd}_{\text{Cat}} \downarrow \text{sm} \) consisting of those morphisms

\[
\begin{array}{ccccccc}
I & \xleftarrow{s} & E & \xrightarrow{p} & B & \xrightarrow{t} & I \\
\downarrow{=} & & \downarrow{pb} & & \downarrow{=} & & \downarrow{=} \\
1 & \xleftarrow{p_S} & P & \xrightarrow{b} & P & \xrightarrow{=} & 1
\end{array}
\]

such that \( I \) is discrete and the functor \( b \) is a discrete fibration.

**Proof.** We first verify that \( \mathcal{N} \) is fully faithful. Given collections \( S \) and \( T \), and a morphism

\[
\begin{array}{ccccccc}
I & \xleftarrow{s_S} & E_S & \xrightarrow{p_S} & B_S & \xrightarrow{t_S} & I \\
\downarrow{=} & & \downarrow{pb} & & \downarrow{=} & & \downarrow{=} \\
J & \xleftarrow{s_T} & E_T & \xrightarrow{p_T} & B_T & \xrightarrow{t_T} & J
\end{array}
\]

\( \mathcal{N} S \to \mathcal{N} T \), one defines \( F : S \to T \) with object map \( f = f_0 \), and with effect on arrows given by the object map of \( f_1 \). Since \( f_1 b_T = b_S \) \( F \)'s arrow map is equivariant. By definition \( \mathcal{N} F = (f_0, f_1, f_2) \), and this equation determines \( F \) uniquely.

For a morphism into \((1, \text{sm})\) in \( \text{PolyEnd}_{\text{Cat}} \) as in the statement, it suffices to exhibit it has \( \mathcal{N} T \) for some collection \( T \). We take the set of objects of \( T \) to be \( I \), and the set of arrows of \( T \) to be \( B \). The target of \( \alpha \in B \) is taken to be \( t \alpha \). Since \((e, b)\) is the structure on \( p \) of a \( U^p \)-fibration, \( p^{-1}\{\alpha\} \) is a finite linearly ordered set, and applying \( s \) componentwise to this produces the source sequence of \( \alpha \) in \( T \). Denoting by \( n \) the length of this sequence and regarding \( \rho \in \Sigma_n \) as an arrow of \( \mathbb{P} \), \( \rho \) lifts to a unique morphism of \( B \) with codomain \( \alpha \) since \( b \) is a discrete fibration, and we denote this unique morphism as \( \rho : \alpha \rho \to \alpha \). Thus we have the required
symmetric group actions, and their functoriality is just that of \( b \). By construction \( \mathcal{N}T \) is the morphism in \( \text{PolyEnd}_{\text{Cat}} \) of the statement.

An operad is a collection \( T \), with object set denoted \( I \), together with

- (units): for \( i \in I \), an arrow \( 1_i : (i) \to i \).
- (compositions): given an arrow \( \alpha : (i_j)_{1 \leq j \leq n} \to i \) of \( T \), and a sequence \( (\beta_j : (i_{jk})_{1 \leq k \leq m_j} \to i_j)_j \) of arrows of \( T \), their composite is an arrow

\[
(\alpha \circ (\beta_j)_j) : (i_{jk})_{1 \leq k \leq m_j} \to i
\]


is the sequence of length \( (m_1 + \ldots + m_n) \) obtained by concatenating the domains of the \( y_j \).

This data must satisfy the following axioms. The unitality and associativity of composition say that given

\[
\alpha : (i_j)_{1 \leq j \leq n} \to i \quad \beta_j : (i_{jk})_{1 \leq k \leq m_j} \to i_j \quad \gamma_{jk} : (i_{jk})_{1 \leq l \leq p_{jk}} \to i_{jk}
\]

one has

\[
1_i \circ (\alpha) = \alpha = \alpha \circ (1_{i_j})_j \quad \alpha \circ (\beta_j \circ (\gamma_{jk})_j)_j = (\alpha \circ (\beta_j)_j) \circ (\gamma_{jk})_j.
\]

Equivariance of composition says that given \( \rho \in \Sigma_n \) and \( \rho_j \in \Sigma_{m_j} \) for \( 1 \leq j \leq n \), one has

\[
(\alpha \circ (\beta_j)_j)(\rho(\rho_j)_j) = (\alpha \rho \circ (\beta_j \rho_j)_j)
\]

where \( \rho(\rho_j)_j \) is the permutation of \( \Sigma_{\sum_{i=1}^n m_i} \) symbols given by permuting the \( n \)-blocks \((m_1, \ldots, m_n)\) using \( \rho \), and permuting the elements within the \( j \)-th block using \( \rho_j \). A morphism \( S \to T \) of operads is a morphism \( F : S \to T \) of the underlying collections, with underlying object map denoted as \( f : I \to J \), such that

\[
F1_i = 1_{fi} \quad F(\alpha \circ (\beta_j)_j) = (F\alpha \circ (F\beta_j)_j)
\]

for all objects \( i \) of \( S \), and arrows \( \alpha \) and \( (\beta_j)_j \) of \( S \) as above. The category of operads and their morphisms is denoted \( \text{Opd} \).

**Proposition 3.2.** The functor \( \mathcal{N} \) lifts to a functor \( \overline{\mathcal{N}} \) making the square

\[
\begin{align*}
\text{Opd} & \xrightarrow{\mathcal{N}} \text{PolyMnd}_{\text{Cat}} & \downarrow \text{sm} \\
\text{Coll} & \xrightarrow{\mathcal{N}} \text{PolyEnd}_{\text{Cat}} & \downarrow \text{sm}
\end{align*}
\]

in which the vertical functors are the forgetful functors, a pullback.

**Proof.** We shall first establish a bijection between operad structures on a collection \( T \), and polynomial monad structures on \( P_T \) making \( \mathcal{N}T \) a morphism of polynomial monads. Second, given collections \( S \) and \( T \) and a morphism \( F : S \to T \) of their underlying collections, we shall prove that \( F \) is a morphism of operads iff \( \mathcal{N}F \) is a morphism of the corresponding polynomial monads over \( \text{sm} \).

Let \( T \) be a collection with object set \( I \). To give units for \( T \) is to give a functor \( u_{T,1} : I \to B_T \) such that: (1) \( t_{1} u_{T,1} = 1_{i} \), (2) for each \( i \) the fibre \( p_{T}^{-1}\{u_{T,1}i\} \) consists of a unique element \( u_{T,2i} \), and (3) \( su_{T,2i} = i \). Since \( B_T \) sends elements with singleton fibres to \( 1 \in \mathbb{P} \) and \( e_T \) sends the unique elements of those fibres to \( (1,1) \in \mathbb{P}_e \), \( \mathcal{N}T \) commutes with these unit maps and those of \( \text{sm} \). Thus to give units for \( T \) is to give \( u_T : 1_1 \to P_T \) with respect to which \( \mathcal{N}T \) is compatible.
We now characterise compositions for an operad structure on \( T \) in similar terms. The polynomial \( P_T \circ P_T \) is formed as

\[
\begin{array}{c}
\text{An object of } B_T^{(2)} \text{ can be identified with a functor } b : [0] \to B_T^{(2)}. \text{ Writing } \alpha = qb, \\
\text{an object of } B_T^{(2)} \text{ may be regarded as the data: (1) an arrow } \alpha : (i_j)_{1 \leq j \leq n} \to i \\
of T \text{ viewed also as a functor } \alpha : [0] \to B_T, \text{ and (2) a functor } b : [0] \to B_T^{(2)} \text{ over } B_T. \text{ The functor } q \text{ is the effect of } \Pi_{p_T} \text{ on the functor } B_T \times I E_T \to E_T, \text{ and the pullback of } \alpha : [0] \to B_T \text{ along } p_T \text{ is the discrete subcategory of } E_T \text{ consisting of the pairs } (\alpha, j) \text{ for } 1 \leq j \leq n. \text{ Thus (2) amounts to giving an arrow } \beta_j \text{ of } T \text{ for each } 1 \leq j \leq n \text{ whose target is } i_j, \text{ and so an object of } B_T^{(2)} \text{ is exactly the data } (\alpha, \beta_j) \text{ that can be composed in the multicategory } T.
\end{array}
\]

Similarly an arrow of \( B_T^{(2)} \) can be identified with a functor \([1] \to B_T^{(2)} \) over \( B_T \). This first datum is just an arrow of \( B_T \), and so is of the form \( \rho : \alpha \rho \to \alpha \), where \( \alpha : (i_j)_{1 \leq j \leq n} \to i \) is an arrow of \( T \), and \( \rho \in \Sigma_n \). Pulling back the functor \([1] \to B_T \) so determined along \( p_T \) produces a category with objects of the form \( (\alpha, j) \) or \( (\alpha, j) \) for \( 1 \leq j \leq n \), and invertible arrows \( (\alpha, j) \to (\alpha, j) \). Thus the second piece of data determining an arrow of \( B_T^{(2)} \) amounts to giving morphisms \( \rho_j : \beta_j \rho_j \to \beta \) in \( B \), such that \( t\beta_j = i_j \) for each \( j \). Thus the general form of an arrow of \( B_T^{(2)} \)

\[
(4) \quad (\rho, (\rho_j)_{1 \leq j \leq n}) : (\alpha \rho, (\beta_j \rho_j)_{j \leq j \leq n}) \to (\alpha, (\beta_j)_{j \leq j \leq n})
\]

where \( \alpha : (i_j)_{1 \leq j \leq n} \to i \) is an arrow of \( T \), \( \rho \in \Sigma_n \), and for each \( j \), \( \beta_j : (i_{jk})_{1 \leq k \leq m_j} \to i_j \) is an arrow of \( T \) and \( \rho_j \in \Sigma_{m_j} \).

A description of \( E_T^{(2)} \) is now easily obtainable, since \( E_T^{(2)} \) is obtained by pulling back the functor \( B_T^{(2)} \to B_T \) which we now know explicitly. So an object of \( E_T^{(2)} \) consists of \( (\alpha, (\beta_j)_{j \leq j \leq n}) \), where \( \alpha : (i_j)_{1 \leq j \leq n} \to i \) and \( \beta_j : (i_{jk})_{1 \leq k \leq m_j} \to i_j \) are arrows of \( T \), \( 1 \leq j \leq n \) and \( 1 \leq k \leq m_j \). Morphisms of \( E_T^{(2)} \) are of the form

\[
(5) \quad (\rho, (\rho_j)_{1 \leq j \leq n}) : (\alpha \rho, (\beta_j \rho_j)_{j \leq j \leq n}) \to (\alpha, (\beta_j)_{j \leq j \leq n})
\]

and the explicit descriptions of the functors \( s_T^{(2)}, p_T^{(2)} \) and \( t_T^{(2)} \) are now self-evident.

Given these details, an object map for a functor \( m_{T,1} : B_T^{(2)} \to B_T \) amounts to assignments \( (\alpha, (\beta_j)_{j \leq j \leq n}) \to \alpha \circ (\beta_j)_{j \leq j \leq n} \). Giving \( m_{T,1} \) on arrows amounts to assigning to (4), an arrow \( (\alpha \rho) \circ (\beta_j \rho_j)_{j \leq j \leq n} \to \alpha \circ (\beta_j)_{j \leq j \leq n} \) of \( B_T \), and the compatibility of \( m_{T,1} \) with \( NT \) and the corresponding component of \( \text{sm} \)’s multiplication, amounts to the underlying permutation of this arrow of \( B_T \) being obtained via the substitution of permutations, which corresponds to equivariance. To say that the target of \( \alpha \circ (\beta_j)_{j \leq j \leq n} \) is that of \( \alpha \) for all \( (\alpha, (\beta_j)_{j \leq j \leq n}) \), is to say that \( m_1 t_T = t_T^{(2)} \). A functor \( m_{T,2} : E_T^{(2)} \to \)
$E_T$ providing the other component of $m_T: P_T \circ P_T \to P_T$ is determined by its restrictions to the fibres of $p_T^{(2)}$ which are finite discrete, and giving these amounts to specifying that for all $(\alpha_i, (\beta_j)_j)$, the domain of the composite $\alpha \circ (\beta_j)_j$ is the concatenation of the domains of the $\beta_j$ as for composition in an operad. In summary, to give $T$ a composition operation is to give $m_T: P_T \circ P_T \to P_T$ in $\text{Poly}_{\text{Cat}}$. The straight forward though tedious verification that the unit and associative laws for $(u_T, m_T)$ correspond with the unitality and associativity of composition for $T$ is left to the reader.

Let $S$ and $T$ be collections and $F: S \to T$ be a morphism of their underlying collections. To say $NF$ is compatible with units amounts to the equation $F_1u_S, 1 = u_T, 1$, the equation $F_2u_S, 2 = u_T, 2$ being a consequence, and this in turn is equivalent to saying that $F$ sends identities in $S$ to identities in $T$. We leave to the reader the straight forward verification that $NF$’s compatibility with multiplications amounts to the formulae $F(\alpha \circ (\beta_j)_j) = F\alpha \circ (F\beta_j)_j$ expressing $F$’s compatibility with composition.

By propositions 3.1 and 3.2 we have

**Theorem 3.3.** The functor $\overline{N}$ restricts to an equivalence between $\text{Opd}$ and the full subcategory of $\text{PolyMnd}_{\text{Cat}} \downarrow \text{sm}$ consisting of those monad morphisms

$$
\begin{array}{c}
\begin{array}{cccc}
I & \xleftarrow{s} & E & \xrightarrow{p} & B & \xrightarrow{t} & I \\
\downarrow & & \downarrow e & & \downarrow pb & & \downarrow b & & \downarrow = \\
1 & \xleftarrow{P} & P & \xrightarrow{U^p} & P & \xrightarrow{U^\tau} & 1
\end{array}
\end{array}
$$

such that $I$ is discrete and the functor $b$ is a discrete fibration.

**Remark 3.4.** A club in the sense of Max Kelly [10, 11] can be identified as a 2-monad $A$ on $\text{Cat}$ together with a cartesian monad morphism $\phi: A \to \text{Sm}$. In general when one has a cartesian monad morphism into a polynomial monad, the domain monad is also easily exhibited as polynomial, and so clubs can be identified as those objects

$$
\begin{array}{c}
\begin{array}{cccc}
I & \xleftarrow{s} & E & \xrightarrow{p} & B & \xrightarrow{t} & I \\
\downarrow & & \downarrow e & & \downarrow pb & & \downarrow b & & \downarrow = \\
1 & \xleftarrow{P} & P & \xrightarrow{U^p} & P & \xrightarrow{U^\tau} & 1
\end{array}
\end{array}
$$

of $\text{PolyMnd}_{\text{Cat}} \downarrow \text{sm}$ such that $I = 1$.

**Remark 3.5.** A $\text{Cat}$-operad is defined in the same way as an operad is, except that the homs are categories and the units and compositions define functors, this being an instance of how the notion of operad can be enriched. Equivalently denoting by $\text{Opd}_I$ the category of operads with objects set $I$ and morphisms whose object function is $1_I$, a $\text{Cat}$-operad with set of objects $I$ is a category internal to $\text{Opd}_I$. As explained in [26] pullbacks in $\text{Poly}_{\text{Cat}}(I, I)$ are formed componentwise, and its straight forward to verify that the restriction

$$
\text{Opd}_I \to \text{Poly}_{\text{Cat}}(I, I)
$$
of \( \mathbf{N} \) preserves pullbacks. From this it is straightforward to see that one can identify \( \mathbf{Cat} \)-operads with objects

\[
\begin{array}{c}
I & \xleftarrow{s} & E & \xrightarrow{p} & B & \xrightarrow{t} & I \\
\downarrow & & \downarrow & & \downarrow & & \downarrow \\
1 & \xleftarrow{\mathbf{P}} & \mathbf{P} & \xrightarrow{\mathbf{U}} & \mathbf{P} \end{array}
\]

of \( \text{PolyMnd}_{\mathbf{Cat}} \downarrow \text{sm} \), together with the structure of a split fibration on \( \mathfrak{b} \). More generally an object of \( \text{PolyMnd}_{\mathbf{Cat}} \downarrow \text{sm} \) in which \( \mathfrak{b} \) is a cloven fibration (not necessarily split) is exactly an operadic category in the sense of section 1.2 of [22].

**Remark 3.6.** A category can be regarded as an operad \( T \) in which the source of every arrow is a sequence of length 1, which is so iff in its underlying polynomial depicted on the left

\[
\begin{array}{c}
I & \xleftarrow{s_T} & E_T & \xrightarrow{p_T} & B_T & \xrightarrow{t_T} & I \\
\downarrow & & \downarrow & & \downarrow & & \downarrow \\
1 & & \mathbf{P}_T & \xrightarrow{\mathbf{U}_T} & \mathbf{P} \end{array}
\]

\( p_T \) is an isomorphism. In this case \( E_T \) and \( B_T \) are discrete, and \( \mathfrak{b}_T : B_T \to \mathbf{P} \) and \( e_T : E_T \to \mathbf{P}_T \) are determined uniquely by the polynomial \( (s_T, p_T, t_T) \). For any polynomial as on the right in the previous display in which \( p \) is an isomorphism and \( E \) and \( B \) are discrete, one as a unique isomorphism \( (s, p, t) \cong (s p^{-1}, 1_B, t) \) with a span of sets. Thus the equivalence of theorem 3.3 essentially restricts to an equivalence of categories which on objects identifies a category with its corresponding monad in \( \text{Span}_{\mathbf{Set}} \).

An operad \( T \) with object set \( I \) determines a 2-monad \((I, P_T)\) in the 2-bicategory \( \text{Poly}_{\mathbf{Cat}} \), and so by means of \( \mathbf{P}_{\mathbf{Cat}} \), a 2-monad on \( \mathbf{Cat}/I \).

**Notation 3.7.** Given an operad \( T \) with object set \( I \), we also denote the associated 2-monad on \( \mathbf{Cat}/I \) as \( T \).

**Example 3.8.** The terminal operad which has one object and a unique arrow of with source of length \( n \), is usually denoted as \( \text{Com} \). Its corresponding polynomial is \( \text{sm} \). Following notation 3.7 one thus has \( \text{Com} = \text{Sm} \).

We now turn to the task of giving an explicit description of this 2-monad \( T \) on \( \mathbf{Cat}/I \). Let \( X \in \mathbf{Cat}/I \). We regard \( X \) both as \( X \to I \) a category equipped with a functor into \( I \), and as \((X_i)_{i \in I}\) an \( I \)-indexed family of categories.

**Definition 3.9.** An *operation of \( T \) labelled in \( X \)* is a pair \( \alpha, (x_j)_j \), where \( \alpha : (i_j)_j \to i \) is an arrow of \( T \), and \( x_j \in X_{i_j} \). A *morphism \( (\alpha, (x_j)_j) \to (\beta, (y_j)_j) \)* is a pair \( (\rho, (\gamma_j)_j) \) where \( \rho \) is a permutation such that \( \alpha = \beta \rho \), and \( \gamma_j : x_j \to y_{\rho(j)} \) is a morphism of \( X_{\rho(j)} \) for each \( j \).

It is also useful to depict a labelled operation \( (\alpha, (x_j)_j) \) of definition 3.9 as

\[
\begin{array}{c}
\circ \\
\downarrow \quad i_i \quad \ldots \quad i_n \\
\circ \\
\downarrow i
\end{array}
\]
and in such diagramatic terms, a morphism amounts to a shuffling of the inputs of the operations, together with a levelwise family of morphisms of $X$. Operations of $T$ labelled by $X$ form a category, and the assignation of codomains of the labelled operations gives a functor from this category into $I$. Applying the general calculation of example 2.2 to the polynomial $(s_T, p_T, t_T)$ one obtains

**Lemma 3.10.** Let $T$ be a collection with object set $I$ and $X \in \text{Cat}/I$. Then $TX$ may be identified with the category of operations of $T$ labelled in $X$.

Similarly one can unpack the explicit description of $Tf : TX \to TY$ given $f : X \to Y$ in $\text{Cat}/I$. By definition $Tf$ is induced from the functoriality of pullbacks and distributivity pullbacks in

![Diagram](6)

and then by tracing through the explicit descriptions as in example 2.2 one can verify

**Lemma 3.11.** Let $T$ be a collection with object set $I$. Given a morphism $f : X \to Y$ in $\text{Cat}/I$, one has

$$Tf(\alpha, (x_j)_j) = (\alpha, (fx_j)_j) \quad Tf(\rho, (\beta_j)_j) = (\rho, (f\beta_j)_j).$$

Remembering that the pullbacks and distributivity pullbacks that appear in (6) enjoy a 2-dimensional universal property, one can in much the same way verify

**Lemma 3.12.** Let $T$ be a collection with object set $I$. Given morphisms $f$ and $g : X \to Y$ in $\text{Cat}/I$ and a 2-cell $\phi : f \to g$, one has

$$(T\phi)_{(\alpha, (x_j)_j)} = (\alpha, (fx_j)_j) \xrightarrow{(1, (\phi x_j)_j)} (\alpha, (gx_j)_j).$$

Lemmas 3.10-3.12 together describe the endo-2-functor of $\text{Cat}/I$ corresponding to a collection $T$ with object set $I$ in terms of labelled operations. We now extend this to a description of the 2-monad corresponding to an operad. In the proof of proposition 3.2 we obtained an explicit understanding of how the identity arrows arrows of an operad provide the unit data for a monad in $\text{Poly}_{\text{Cat}}$. Putting this together with the explicit description of the homomorphism $P_{\text{Cat}} : \text{Poly}_{\text{Cat}} \to \text{2-Cat}$ of 2-bicategories, we obtain

**Lemma 3.13.** Let $T$ be an operad with object set $I$. Given an object $X$ of $\text{Cat}/I$ one has

$$\eta^T_X(x) = (1_i, (x)) \quad \eta^T_X(\beta) = (1_i, (\beta))$$

for any $x$ and $\beta : x \to y$ in $X_i$.  

Similarly from the explicit understanding of how the compositions of an operad give rise to the multiplication data for a monad in $\textbf{Poly}_{\text{Cat}}$ obtained in proposition 3.2, we further obtain

**Lemma 3.14.** Let $T$ be an operad with object set $I$ and $X \in \text{Cat}/I$. Then the effect of $\mu^T_X$ on objects is given by

$$\mu^T_X(\alpha,(\alpha_j,(x_{jk})_j)_j) = (\alpha(\alpha_j)_j,(x_{jk})_j)_j.$$  

The effect of $\mu^T_X$ on an arrow

$$(\rho,(\rho_j, (\beta_{jk})_j)_j) : (\alpha\rho,(\alpha_j\rho_j,(x_{jk})_j)_j)_j \rightarrow (\alpha,(\alpha_j,(y_{jk})_j)_j)_j$$

of $T^2X$ is

$$(\rho(\rho_j)_j,(\beta_{jk})_j)_j : ((\alpha(\alpha_j)_j)(\rho(\rho_j)_j),(x_{jk})_j)_j \rightarrow (\alpha(\alpha_j)_j,(y_{jk})_j)_j).$$

In diagramatic terms the object map of $\mu^T_X$ may be depicted as

![Diagram](image)

4. Categorical algebras of operads as weak operad morphisms into $\text{Cat}$

As we have seen, an operad $T$ with set of objects $I$ can be regarded as a 2-monad on $\text{Cat}/I$. In this section we see that lax, colax, pseudo and strict algebras for $T$ correspond respectively to lax, colax, pseudo and strict morphisms of $\text{Cat}$-operads into $\text{Cat}$ viewed as a $\text{Cat}$-operad. In the case where $T$ is an operad with only arrows of arity 1, so that $T$ is itself just a category, this generalises the 2-monads described in [3] whose algebras are (the appropriately weakened) functors $T \rightarrow \text{Cat}$.

We denote by $\text{Cat}^{\text{op}}$ the $\text{Cat}$-operad whose objects are small categories and whose homs are given by the functor categories $\text{Cat}((A_j)_{1 \leq j \leq n}; B) = \prod_{j=1}^{n} A_j, B$.

**Definition 4.1.** Let $T$ be an (unenriched) operad with object set $I$. A lax morphism of operads $H : T \rightarrow \text{Cat}$ consists of

- $\forall i \in I$, a category $H_i$,
- $\forall \alpha : (i_j)_{1 \leq j \leq n} \rightarrow i$ in $T$, a functor $H_\alpha : \prod_{j=1}^{n} H_{i_j} \rightarrow H_i$,
- $\forall \alpha$ and $\rho \in \Sigma_n$, a natural transformation $\xi_{\alpha, \rho} : H_{\alpha\rho}c_\rho \rightarrow H_\alpha$, where $c_\rho : \prod_{j=1}^{n} H_i = \prod_{j=1}^{n} H_{i_{\rho_j}}$ is given by permuting the factors according to $\rho$.
- $\forall i \in I$, a natural transformation $\nu_i : 1_{H_i} \rightarrow H_i$,
- $\forall (\alpha, (\beta_j)_j)$ where $\alpha : (i_j)_{1 \leq j \leq n} \rightarrow i$ and $\beta_j : (i_{jk})_k \rightarrow i$ are in $T$, a natural transformation $\sigma_{\alpha, (\beta_j)_j} : H_\alpha(\prod_{j=1}^{n} H_{\beta_j}) \rightarrow H_{\alpha(\beta_j)_j}$.
such that $\xi_{\alpha,1} = \text{id}$, $\xi_{\alpha,\rho_1 \rho_2} = \xi_{\alpha,\rho_1}(\xi_{\alpha_1 \rho_1,\rho_2})$, and

\[
H_\alpha \xrightarrow{\text{id} \prod_j \nu} H_\alpha \prod_j H_{1_{ij}} \xrightarrow{} H_\alpha \xrightarrow{\nu \text{id}} H_1 H_\alpha
\]

\[
H_\alpha \xrightarrow{H_{\alpha \rho_1 \rho_2}} H_\alpha \xrightarrow{H_{\alpha(\rho_1)\rho_2}} H_{\alpha(\rho_1)(\rho_2)}
\]

\[
H_\alpha \prod_j H_{\beta_j \beta_j} \prod_{j,k} H_{\gamma_{jk}} \xrightarrow{\text{id} \cdot \text{sh}_n} H_\alpha \prod_j (H_{\beta_j} \prod_k H_{\gamma_{jk}})
\]

\[
H_\alpha \prod_{j,k} H_{\gamma_{jk}} \xrightarrow{\sigma \cdot \text{id}} H_\alpha \prod_j H_{\alpha(\beta_{jk})(\gamma_{jk})} \xrightarrow{\sigma} H_\alpha \prod_{j,k} H_{\alpha(\beta_{jk})(\gamma_{jk})}
\]

\[
H_\alpha \prod_{j,k} H_{\gamma_{jk}} \xrightarrow{\sigma \cdot \text{id}} H_\alpha \prod_{j,k} H_{\alpha(\beta_{jk})(\gamma_{jk})}
\]

commute, where $\text{sh}_n \in \Sigma_{2n}$ is the “shuffle” permutation$^2$.

**Definition 4.2.** In the context of definition 4.1 the functors $H_\alpha$ are called the *products*, and the natural transformations $\xi_{\alpha,\rho}$, $\nu_i$ and $\sigma_{\alpha,\beta_{ij}}$ are called the *symmetries*, *units* and *substitutions* for $H$. When the units and substitutions are invertible, $H$ is said to be a *pseudo morphism*, and when they are identities $H$ is said to be a *strict morphism*.

Clearly the symmetries are isomorphisms in general. However even for a strict morphism, the symmetries need not be identities.

**Definition 4.3.** In the context of definitions 4.1 and 4.2, a lax morphism $H : T \to \textbf{Cat}$ is commutative when its symmetries are identities.

For any operad $T$ with object set $I$ and symmetric monoidal category $(\mathcal{V}, \otimes)$, it is standard to consider algebras of $T$ in $\mathcal{V}$. In explicit terms such an algebra consists of an object $H_i$ in $\mathcal{V}$ for each $i \in I$, morphisms $H_\alpha : \bigotimes_j H_{i_j} \to H_i$ for any $\alpha : (i_j)_j \to i$ in $T$, satisfying

\[
H_{\alpha \rho} = H_\alpha c_\rho \\
H_{1_{ij}} = 1_{H_i} \\
H_{\alpha(\beta_{ij})} = H_\alpha(\bigotimes_j H_{\beta_{ij}})
\]

for all $i \in I$, $\alpha : (i_{1j})_{1 \leq j \leq n} \to i$ and $\beta_{ij} : (i_{jk})_k \to i_j$ in $T$, and $\rho \in \Sigma_n$. Thus in particular one has

**Example 4.4.** A commutative strict morphism $H : T \to \textbf{Cat}$ is the same thing as an algebra of the operad $T$ in $(\textbf{Cat}, \times)$.

On the other hand the most fundamental general class of examples which are rarely commutative is

**Example 4.5.** A symmetric monoidal category $(\mathcal{V}, \otimes)$ is the same thing as a pseudo morphism $\mathcal{V} : \text{Com} \to \textbf{Cat}$. The unique object of $\text{Com}$ is sent to the underlying category also denoted as $\mathcal{V}$, if $\alpha$ is the unique morphism of $\text{Com}$ of arity $n$ then $\mathcal{V}_\alpha$ is the tensor product functor $\otimes : \mathcal{V}^n \to \mathcal{V}$, and the rest of the data of $\mathcal{V} : \text{Com} \to \textbf{Cat}$ corresponds exactly to the coherence morphisms of $\mathcal{V}$. Lax morphisms $\text{Com} \to \textbf{Cat}$ are also very well studied and are referred to either as symmetric lax monoidal categories or functor operads in the literature.

$^2$As an endofunction of $\{1, \ldots, 2n\}$, $\text{sh}_n$ is defined by $\text{sh}_n(j) = \frac{j + 1}{2}$ when $j$ is odd, and $\text{sh}_n(j) = n + \frac{j}{2}$ when $j$ is even.
It is also interesting to consider the following generalisation of example 4.5
definable for any operad \( T \).

**Example 4.6.** Let \( \mathcal{V} \) be a symmetric lax monoidal category and \( T \) be an operad. Then one has a lax morphism \( \mathcal{V}_* : T \to \mathbf{Cat} \) in which \( (\mathcal{V}_*)_n = \mathcal{V} \), \( (\mathcal{V}_*)_a \) where \( a \)’s source is a sequence of length of \( n \) is \( \otimes : \mathcal{V}^n \to \mathcal{V} \), and the units, substitutions and symmetries are given by the coherences for \( \mathcal{V} \)’s symmetric lax monoidal structure. When \( \mathcal{V} \) is a symmetric monoidal category, \( \mathcal{V}_* \) is a pseudo morphism \( T \to \mathbf{Cat} \). In particular when \( \mathcal{V} \) is the terminal category, \( \mathcal{V}_* \) is denoted as \( 1 : T \to \mathbf{Cat} \), which clearly corresponds via example 4.4 to the terminal algebra in \((\mathbf{Cat}, \times)\) of the operad \( T \).

**Example 4.7.** As discussed in remark 3.6 an operad \( T \) with only unary arrows is a category. In this case a lax, pseudo or strict morphism \( T \to \mathbf{Cat} \) is the same thing as a lax functor, pseudo functor or functor \( T \to \mathbf{Cat} \), in the usual senses, respectively. Moreover all such morphisms \( T \to \mathbf{Cat} \) are commutative.

The various types of morphisms \( H : T \to \mathbf{Cat} \) defined above will now be exhibited as being structure on the underlying object of \( \mathbf{Cat}/I \) whose fibre over \( i \in I \) is the category \( H_i \). We abuse notation and refer to this object of \( \mathbf{Cat}/I \) also as \( H \), or as \( H \to I \).

For a general 2-monad \((T, \eta, \mu)\) on a 2-category \( \mathcal{K} \), we recall now the various types of \( T \)-algebra structure one can have on an object \( A \in \mathcal{K} \). A lax \( T \)-algebra structure on \( A \) consists of an arrow \( a : TA \to A \), coherence 2-cells \( \pi_0 : 1_A \to a_\eta A \) and \( \pi_2 : aT(a) \to a\mu A \) such that

\[
\begin{array}{ccc}
T(a) & \longrightarrow & T(a)T^2(a) \\
\pi_0 \downarrow & & \downarrow \pi_2 \\
\pi_0 a & \longrightarrow & a\mu A \\
\end{array}
\]

\[
\begin{array}{ccc}
a & \longrightarrow & aT(a) \\
\pi_2 \downarrow & & \downarrow \pi_0 \\
aT(\pi_0) & \longrightarrow & \pi_0 a \\
\end{array}
\]

\[
\begin{array}{ccc}
a T(a) & \longrightarrow & aT(\eta A) \\
\pi_0 \downarrow & & \downarrow \pi_2 \\
\pi_0 aT(\pi_2) & \longrightarrow & a\mu A \\
\end{array}
\]

\[
\begin{array}{ccc}
a T(a)T(\mu A) & \longrightarrow & aT(\eta A) \\
\pi_2 \downarrow & & \downarrow \pi_0 \\
aT(\pi_2) & \longrightarrow & \pi_0 a \\
\end{array}
\]

\[
\begin{array}{ccc}
a T(\eta A) & \longrightarrow & aT(a) \\
\pi_0 \downarrow & & \downarrow \pi_2 \\
\pi_0 aT(\pi_0) & \longrightarrow & a\mu A \\
\end{array}
\]

\[
\begin{array}{ccc}
a & \longrightarrow & aT(a) \\
\pi_2 \downarrow & & \downarrow \pi_0 \\
aT(\pi_0) & \longrightarrow & \pi_0 a \\
\end{array}
\]

commute. We denote a lax \( T \)-algebra as a pair \((A, a)\) leaving the coherence data \( \pi_0 \) and \( \pi_2 \) implicit. When these coherences are isomorphisms \((A, a)\) is called a *pseudo* \( T \)-algebra, and when they are identities \((A, a)\) is called a *strict* \( T \)-algebra.

**Theorem 4.8.** Let \( T \) be an operad with object set \( I \), regard it as a 2-monad on \( \mathbf{Cat}/I \) as in section 3, and let \( H \to I \) be an object of \( \mathbf{Cat}/I \). To give \( H \) the structure of a lax, pseudo or strict \( T \)-algebra is to give it the structure of a lax, pseudo or strict morphism \( H : T \to \mathbf{Cat} \), respectively, in the sense defined above.

**Proof.** Before proceeding to unpack what the action \( a : TH \to H \) amounts to, we observe first that there is a canonical factorisation system\(^3\) on \( TH \). Recall that a general arrow of \( TH \) is of the form \((\rho, (\gamma_j)_j) : (\alpha_\rho, (x_j)_j) \to (\alpha, (y_j)_j)\) where \( \alpha : (i_j)_{1 \leq j \leq n} \to i \) is from \( T \), \( \rho \in \Sigma_n \) and \( \gamma_j : x_j \to y_{j\rho} \) is from \( H_{\rho j} \), and we say that \((\rho, (\gamma_j)_j)\) is *levelwise* when \( \rho \) is an identity, and *permutative* when the \( \gamma_j \) are all identities. Both these types of maps are closed under composition and contain the identities, a map is an identity iff it is both levelwise and permutative, and a

\(^3\)This is in a very strict algebraic sense, of being a decomposition of \( TH \) as a composite monad in \( \mathbf{Span}_{\mathbf{Set}} \) via a distributive law, a situation that was studied in (ref. to Rosebrugh-Wood).
general map \((\rho, (\gamma_j)_j)\) factors in a unique way

\[
(\alpha, (x_j)_j) \overset{(1_a, (\gamma_j)_j)}\longrightarrow (\alpha, (y_j)_j) \overset{(\rho, (1_{y_j})_j)}\longrightarrow (\alpha, (y_j)_j)
\]

as a levelwise map followed by a permutative map. Denoting the subcategory of \(TH\) containing all the levelwise maps as \(T_L H\), and the subcategory of \(TH\) containing all the permutative maps as \(T_R H\), to give a functor \(f : TH \to C\) into any category \(C\), is to give functors \(f_L : T_L H \to C\) and \(f_R : T_R H \to C\) which agree on objects, and with arrow maps compatible in the following sense – if \((r : a \to b, l : b \to c)\) is a composable pair in \(TH\) in which \(l\) is levelwise and \(r\) permutative, then \(f_L(l)f_R(r) = f_R(r')f_L(l')\) in \(C\), where \(lr = r'l'\) is the levelwise-permutative factorisation of \(lr\).

The object map of \(a : TH \to H\) gives \(a(\alpha, (y_j)_j) \in H_i\) for each \((\alpha, (y_j)_j)\) as above. Allowing the \((y_j)_j\) to vary, this amounts to the object map of a functor \(H_\alpha : \prod_{j=1}^n H_i \to H_i\) for each \(\alpha\). The arrow map of \(a_L\) gives \(H_\alpha(\gamma_j)_j : H_\alpha(x_j)_j \to H_\alpha(y_j)_j\) in \(H_i\), for each \((\alpha, (y_j)_j)\) as above, and \(\gamma_j : x_j \to y_j\) in \(H_j\). Allowing the \((\gamma_j)_j\) to vary, we see that to give \(a_L\) is to give functors \(H_\alpha\) for all \(\alpha\). The data of the arrow map of \(a_R\) gives, for each \((\alpha, (y_j)_j)\) as above and \(\rho \in \Sigma_n\), a morphism \(H_{\alpha \rho}(y_j)_j \to H_\alpha(y_j)_j\), and allowing the \((y_j)_j\) to vary, this amounts to the components of a natural transformation \(\xi_{\alpha, \rho} : H_{\alpha \rho} \to H_{\alpha c\rho}\). In these terms the functoriality of \(a_R\) amounts to the equations \(\xi_{\alpha, 1} = \text{id}\) and \(\xi_{\alpha, \rho_1 \rho_2} = \xi_{\alpha, \rho_1}(\xi_{\alpha, \rho_1 \rho_2}, c\rho_2)\) and the compatibility of the arrow maps of \(a_L\) and \(a_R\) amounts to the naturality of the \(\xi_{\alpha, \rho}\). Thus we have verified that to give \(a : TH \to H\) is to give the products and symmetries of a lax morphism \(H : T \to \mathbf{Cat}\).

By the explicit description of the unit of the 2-monad \(T\) given in lemma 3.13, to give \(\overline{\alpha}_0 : 1 \to a\eta^T\) is to give \(\nu_i : 1_{H_i} \to H_{1_i}\), and \(\overline{\sigma}_0\) is invertible or an identity iff the \(\nu_i\) are so. By the explicit description of the multiplication of the 2-monad \(T\) given in lemma 3.14, one can similarly reconcile the data of \(\overline{\tau}_2\) with that of the substitutions \(\sigma_{\alpha, (\beta)_j}\). The naturality of \(\overline{\tau}_2\) with respect to levelwise maps corresponds to the naturality of the \(\sigma_{\alpha, (\beta)_j}\), the naturality of \(\overline{\tau}_2\) with respect to permutative maps corresponds to the axiom of compatibility between substitutions and symmetries, and the axioms involving just units and substitutions amount to the lax algebra coherence axioms for \((\overline{\alpha}_0, \overline{\tau}_2)\).

This last result gives an explicit characterisation of the various kinds of algebras of the 2-monad on \(\mathbf{Cat}/I\) associated to an operad \(T\) with object set \(I\). We now proceed to do the same for morphisms of \(T\)-algebras and algebra 2-cells.

**Definition 4.9.** Let \(T\) be an operad with object set \(I\). Suppose that \(H\) and \(K\) are lax morphisms of operads \(T \to \mathbf{Cat}\). A lax-natural transformation

\[(f, \overline{f}) : H \to K\]

consists of

- \(\forall i \in I\), a functor \(f_i : H_i \to K_i\).
- \(\forall \alpha : (i_j)_{1 \leq j \leq n} \to i\) in \(T\), a natural transformation

\[
\overline{f}_\alpha : K_\alpha(\prod_j f_{i_j}) \to f_i H_\alpha
\]
such that

\[
\begin{array}{cccccc}
K_{\alpha \rho \epsilon \rho} \prod_j f_{ij} & \xrightarrow{\overline{\epsilon} \cdot \id} & f_i H_{\alpha \rho \epsilon \rho} & \xrightarrow{\id \cdot \overline{\epsilon} \cdot \id} & \prod_j f_{ij} H_{\alpha \rho \epsilon \rho} & \xrightarrow{\id \cdot \overline{\epsilon} \cdot \id} & \prod_j f_{ij} H_{\alpha \rho \epsilon \rho} \\
\xi^K \cdot \id & \xrightarrow{\id \cdot \xi^H} & \id \cdot \xi^H & \xrightarrow{\id \cdot \nu^H} & \id \cdot \nu^H & \xrightarrow{\id \cdot \nu^H} & \id \cdot \nu^H \\
K_{\alpha \rho \epsilon \rho} \prod_j f_{ij} & \xrightarrow{\overline{\epsilon} \cdot \id} & f_i H_{\alpha \rho \epsilon \rho} & \xrightarrow{\id \cdot \overline{\epsilon} \cdot \id} & \prod_j f_{ij} H_{\alpha \rho \epsilon \rho} & \xrightarrow{\id \cdot \overline{\epsilon} \cdot \id} & \prod_j f_{ij} H_{\alpha \rho \epsilon \rho}
\end{array}
\]

commute. Modifying this definition by reversing the direction of the \( \overline{\epsilon} \) gives the definition of a colax-natural transformation. When the \( \overline{\epsilon} \) are invertible, \( f \) is said to be pseudo-natural, and when they are identities \( f \) is just called a natural transformation.

**Example 4.10.** If as in example 4.7 \( T \) is a category, so that \( H \) and \( K \) can be regarded as lax functors \( T \to \text{Cat} \) in the usual sense, then definition 4.9 in this case gives the usual notions of lax-natural, colax-natural, pseudo-natural and natural transformation \( H \to K \).

**Example 4.11.** Let \( T \) be an operad and \( V \) be a symmetric monoidal category. Then an algebra of the operad \( T \) in \( V \) (recalled just after definition 4.3) is the same thing as a lax natural transformation \( 1 \to V_\bullet \), where \( 1 \) and \( V_\bullet \) are the morphisms \( T \to \text{Cat} \) defined in example 4.6. In the same way one can define algebras of \( T \) within symmetric lax monoidal categories.

In the context of definition 4.9 we write \( f : H \to K \) for the morphism of \( \text{Cat} / I \) whose morphism between fibres over \( i \in I \) is \( f_i \). The notions just defined will now be seen as structure on \( f \).

For a general 2-monad \((T, \eta, \mu)\) on a 2-category \( K \), we recall now the various types of \( T \)-algebra morphism structure one can have on a morphism \( f : A \to B \) in \( K \), where \( A \) and \( B \) underlie lax \( T \)-algebras \((A, a)\) and \((B, b)\). A lax morphism \((A, a) \to (B, b)\) is a pair \((f, \overline{\epsilon})\), where \( f : A \to B \) and \( \overline{\epsilon} : bT(f) \to fa \) such that

![Diagram](diagram.png)

commute. Modifying this definition by reversing the direction of \( \overline{\epsilon} \) gives the notion of colax morphism, when \( \overline{\epsilon} \) is an isomorphism \( f \) is said to be a pseudo morphism, and when \( \overline{\epsilon} \) is an identity \( f \) is said to be a strict morphism.

**Theorem 4.12.** Let \( T \) be an operad with object set \( I \), regard it as a 2-monad on \( \text{Cat} / I \) as in section 3, and let \( f : H \to K \) be a morphism of \( \text{Cat} / I \). Suppose also that one has the structure of lax morphism \( T \to \text{Cat} \) on both \( H \) and \( K \). Then to give \( f \) the structure of lax, colax, pseudo or strict \( T \)-algebra morphism, is to give \( f \)
the structure of lax-natural, colax-natural, pseudo-natural or natural transformation respectively, with respect to the corresponding lax $T$-algebra structures on $H$ and $K$.

**Proof.** Let us write $(a, \pi_0, \pi_2)$ for the action $a : TH \to H$ and coherence cells for the lax $T$-algebra structure on $H$ which corresponds to its lax morphism structure by theorem 4.8, and similarly we write $(b, \pi_0, \pi_2)$ in the case of $K$. One has a component of the lax $T$-algebra coherence datum $\mathcal{F} : bT(f) \to fa$ for each object of $TH$, which is a pair $((\alpha, (y_j)_j))$ where $\alpha : (i_j)_{1 \leq j \leq n} \to i$ is from $T$ and $y_j \in H_{i_j}$, and so such a component is a morphism $K_\alpha(fy_j)_j \Rightarrow fH_\alpha(y_j)_j$. Thus the components of $\mathcal{F} : bT(f) \to fa$ are exactly the components of $\mathcal{F}_\alpha$ for all $\alpha$ for a corresponding lax-natural transformation structure on $f$, and $\mathcal{F}$ is an isomorphism or an identity iff the $\mathcal{F}_\alpha$ are so. The naturality of $\mathcal{F}$ with respect to the levelwise maps in $TH$ corresponds exactly to the naturality of the $\mathcal{F}_\alpha$ for all $\alpha$, and the naturality of $\mathcal{F}$ with respect to the permutative maps corresponds exactly to the axiom on the $\mathcal{F}_\alpha$’s relating to the symmetries of $H$ and $K$. The remaining axioms relating the $\mathcal{F}_\alpha$’s with the units and substitutions of $H$ and $K$ correspond to the lax algebra coherence axioms. \qed

**Remark 4.13.** Reversing the direction of the units and substitutions in the definition of “lax morphism” of operads $H : T \to \mathbf{Cat}$, one obtains the definition of a *colax morphism* $H$. Similarly given a 2-monad $T$ on a 2-category $K$, reversing the direction of the coherence 2-cells $\pi_0$ and $\pi_2$ in the definition of a “lax $T$-algebra structure” on $A \in K$, one obtains the definition of *colax $T$-algebra structure* on $A \in K$. For an operad $T$ with object set $I$, to give $h : H \to I$ the structure of a colax morphism $T \to \mathbf{Cat}$ is to give $h^{op} : H^{op} \to I^{op} = I$ the structure of a lax morphism $T \to \mathbf{Cat}$. To give $h$ the structure of a colax $T$-algebra is to give $h^{op}$ the structure of a lax $T$-algebra. Thus one has versions of theorems 4.8 and 4.12 characterising colax $T$-algebras and various kinds of morphisms between them.

Finally we characterise the algebra 2-cells for the 2-monad on $\mathbf{Cat}/I$ associated to an operad $T$ with object set $I$.

**Definition 4.14.** Let $T$ be an operad with object set $I$. Suppose that $H$ and $K$ are lax morphisms of operads $T \to \mathbf{Cat}$, and that $(f, \mathcal{F})$ and $(g, \mathcal{G}) : H \to K$ are lax-natural transformations. Then a modification $\psi : (f, \mathcal{F}) \Rightarrow (g, \mathcal{G})$ consists of natural transformations $\psi_i : f_i \Rightarrow g_i$ for all $i \in I$, such that

\[
\begin{array}{ccc}
K_\alpha(\prod_j f_{i_j}) & \xrightarrow{\mathcal{F}_\alpha} & f_iH_\alpha \\
\text{id} \cdot \prod_j \psi_i & \downarrow & \psi_i \cdot \text{id} \\
K_\alpha(\prod_j g_{i_j}) & \xrightarrow{\pi_\alpha} & g_iH_\alpha
\end{array}
\]

commutes for all $\alpha : (i_j)_{1 \leq j \leq n} \to i$ in $T$.

In the evident way by reversing the directions of the appropriate coherence cells, one defines modifications between colax natural transformations, and algebra 2-cells between colax morphisms of $T$-algebras.

Recall that for a general 2-monad $T$ on a 2-category $K$, lax $T$-algebras $(A, a)$ and $(B, b)$, and lax morphisms of $T$-algebras $(f, \mathcal{F})$ and $(g, \mathcal{G}) : (A, a) \to (B, b)$, that an *algebra 2-cell* $(f, \mathcal{F}) \Rightarrow (g, \mathcal{G})$ is a 2-cell $\psi : f \Rightarrow g$ in $K$ such that $(\psi a)\mathcal{F} = \mathcal{G}(bT(\psi))$.  

It is straightforward to verify that the modifications of definition 4.14 match up with the algebra 2-cells of $T$ as follows.

**Theorem 4.15.** Let $T$ be an operad with object set $I$, regard it as a 2-monad on $\text{Cat}/I$. Suppose that one has lax morphisms of operads $H$ and $K : T \to \text{Cat}$, and lax-natural (resp. colax-natural) transformations $(f, \overline{f})$ and $(g, \overline{g}) : H \to K$. Then to give a modification $(f, \overline{f}) \to (g, \overline{g})$ is to give an algebra 2-cell between the corresponding lax (resp. colax) morphisms of $T$-algebras.

The various notions of algebra and algebra morphism for a general 2-monad $T$ form the various 2-categories of algebras of $T$, the standard notation for which is recalled partially in

| Name      | Objects         | Arrows         |
|-----------|-----------------|----------------|
| Lax-$T$-Alg | lax $T$-algebras | lax morphisms  |
| Ps-$T$-Alg$_0$ | pseudo $T$-algebras | lax morphisms  |
| Ps-$T$-Alg | pseudo $T$-algebras | pseudo morphisms |
| Ps-$T$-Alg$_s$ | pseudo $T$-algebras | strict morphisms |
| $T$-Alg$_1$ | strict $T$-algebras | lax morphisms  |
| $T$-Alg | strict $T$-algebras | pseudo morphisms |
| $T$-Alg$_s$ | strict $T$-algebras | strict morphisms |

In each case, the 2-cells are just the $T$-algebra 2-cells between the appropriate $T$-algebra morphisms. For the 2-monad $T$ corresponding to an operad it is straightforward to reconcile the compositions one has in these various 2-categories of $T$-algebras with the evident compositions of (co)lax-natural transformations and modifications defined between the various kinds of morphisms $T \to \text{Cat}$.

In particular in the context of example 4.11, the category $\text{Ps-}T$-$\text{Alg}_0(1, V_\bullet)$ can be identified with the usual category of algebras and algebra morphisms of $T$ in $\mathcal{V}$. Thus for each flavour of operad exhibited as arising from a polynomial monad in $[I]$, and each symmetric monoidal category $\mathcal{V}$, for the appropriate choice of $T$, $\text{Ps-}T$-$\text{Alg}_0(1, V_\bullet)$ is the category of that kind of operad in $\mathcal{V}$. For instance, cyclic and modular operads in a symmetric monoidal category are described in this way.

### 5. Commutative operad morphisms into $\text{Cat}$

In this section we exhibit a 2-monad $T/\Sigma$ on $\text{Cat}/I$, whose lax, colax, pseudo and strict algebras are exactly the commutative lax, colax, pseudo and strict morphisms $H$. So it is the strict-$(T/\Sigma)$-algebras that correspond to ordinary morphisms of operads $T \to \text{Cat}$. As explained in the introduction, we will obtain $T/\Sigma$ from $T$ by exhibiting a canonical 2-cell $\alpha_T$

$$ T^1_{\Sigma} \xrightarrow{\delta_T} T \xrightarrow{q_T} T/\Sigma $$

(7)

in the 2-category $\text{Mnd}(\text{Cat}/I)$ of 2-monads on $\text{Cat}/I$ below in definition 5.6, and then taking the universal 1-cell $q_T$ such that $q_T \circ \alpha_T$ is an identity. In the language of 2-category theory, $q_T$ is the coidentifier of $\alpha_T$ in $\text{Mnd}(\text{Cat}/I)$. From this abstract description we will see that $T/\Sigma$ has the required algebras. Recall from the proof of theorem 4.8 that for $H \to I$ in $\text{Cat}/I$, $TH$ has a factorisation system in which
the left class are the levelwise maps and the right class are the permutative maps. The universal property of \( q_T \) means that it is the universal monad morphism which on components sends these permutative maps to identities. In other words at the syntactic level, the monad morphism \( q_T \) is the process of modding out by the symmetric group actions of the operad \( T \).

Before providing \( T_k[1] \Sigma \) and \( \alpha_T \) for (7) some preliminary remarks are in order. To begin with, it turns out that for us, the 2-functor \((-)^{[1]} : \text{Cat} \to \text{Cat}\) which takes any category to its category of arrows, preserves enough distributivity pullbacks. A general result in this direction is lemma 5.1 below. Recall that given functors \( F, G : A \to B \) that a natural transformation \( \phi \) is cartesian when its naturality squares in \( B \) are pullbacks. Given an arrow \( f \) of \( A \), we say that \( \phi \) is cartesian with respect to \( f \) when its naturality square associated to \( f \) is a pullback square, so that \( \phi \) is cartesian iff it is cartesian with respect to all the morphisms of \( A \).

**Lemma 5.1.** A functor \( R : E \to F \) between categories with pullbacks with a pullback preserving left adjoint \( L \), preserves any distributivity pullbacks

\[
\begin{array}{c}
P \ar[r]^p & A \ar[r]^g & B \\
q \ar[u] & \ar[d]^f & \\
Q \ar[r]^r & C
\end{array}
\]

such that the counit of \( L \dashv R \) is cartesian with respect to \( f \).

**Proof.** We denote by \( \varepsilon : LR \to 1_E \) the counit of the adjunction \( L \dashv R \). Given a pullback \((p_2, q_2, r_2)\) around \((Rf, Rg)\) as on the left

we must exhibit \((\gamma, \delta)\) as shown in \( F \) unique making the diagram commute. To this end one has the solid parts of the diagram on the right in \( E \) in the previous display, and the square with vertices \((LX, B, C, LY)\) can be decomposed as

\[
\begin{array}{c}
LX \ar[r]^{Lp_2} & LRA \ar[r]^{LRg} & LRB \ar[r]^{\varepsilon_B} & B \\

LY \ar[r]^{Lq_2} & LR A \ar[r]^{LRf} & LRC \ar[r]^{\varepsilon_C} & C
\end{array}
\]

and so is a pullback by the hypotheses on \( L \) and \( f \). Thus one has \((\alpha, \beta)\) unique as shown making the diagram on the right commute, but to give such \((\alpha, \beta)\) is, by the adjointness \( L \dashv R \), to give the required \((\gamma, \delta)\). \(\square\)
Remark 5.2. To obtain the 2-categorical analogue of lemma 5.1, in which \( L \dashv R \) is a 2-adjunction, one uses the 2-dimensional universal property of distributivity pullbacks in this setting and the 2-dimensional aspects of the adjointness \( L \dashv R \), to adapt the above proof.

Lemma 5.3. If the pullback on the left

\[
\begin{array}{ccc}
P & \xrightarrow{p} & A \xrightarrow{g} B \\
\downarrow{q} & & \downarrow{f} \\
Q & \xrightarrow{r} & C
\end{array}
\]

in \( \text{Cat} \) is a distributivity pullback around \((f, g)\) and \(f\) is a discrete fibration and a discrete opfibration, then the pullback on the right in the previous display is a distributivity pullback around \((f[1], g[1])\).

Proof. The functor \((-) \times [1] : \text{Cat} \to \text{Cat}\) preserves pullbacks since it can be written as the composite

\[
\text{Cat} \xrightarrow{\Delta[1]} \text{Cat}/[1] \xrightarrow{\Sigma_2[1]} \text{Cat}.
\]

The component of the counit of \((-) \times [1] \dashv (-)[1]\) at \(X \in \text{Cat}\) is given by the evaluation functor \(ev_{[1], X} : X[1] \times [1] \to X\). Thus by lemma 5.1, it suffices to show that if \(f : B \to C\) is a discrete fibration and a discrete opfibration, then the naturality square

\[
\begin{array}{ccc}
B[1] \times [1] & \xrightarrow{ev_{[1], B}} & B \\
\downarrow{f[1] \times 1_{[1]}} & & \downarrow{f} \\
C[1] \times [1] & \xrightarrow{ev_{[1], C}} & C
\end{array}
\]

is a pullback. To say this is a pullback on objects is to say that for \(\alpha : c_0 \to c_1\) in \(C\), \(i \in \{0, 1\}\) and \(b \in B\) such that \(fb = c_i\), then there exists a unique \(\beta : b_0 \to b_1\) of \(B\) such that \(f\beta = \alpha\) and \(b_1 = b\). When \(i = 0\), this is exactly the condition that \(f\) be a discrete opfibration, and when \(i = 1\) this corresponds to \(f\) being a discrete fibration. In other words \(f\) is a discrete fibration and a discrete opfibration iff (8) is a pullback on objects.

An arrow \(S\) of \(C[1]\) is a commutative square in the category \(C\), and the category \([1]\) is just the ordinal \(\{0 < 1\}\), and so it has three arrows \(1_0, 1_1\) and the unique arrow \(0 \to 1\). Writing

\[
\begin{array}{ccc}
c_0 & \xrightarrow{\alpha_0} & c_1 \\
\downarrow{\alpha_2} & & \downarrow{\alpha_3} \\
c_2 & \xrightarrow{\alpha_1} & c_3
\end{array}
\]

for the arrow \(S\), regarding it as an arrow \(\alpha_0 \to \alpha_1\) in \(C[1]\), the evaluation functor \(ev_{[1], C}\) acts on arrows by sending \((S, 1_0), (S, 1_1)\) and \((S, 0 \to 1)\) to \(\alpha_2, \alpha_3\) and the diagonal \(c_0 \to c_3\) respectively. To say that (8) is a pullback on arrows is to say that given a square \(S\) in \(C\), an arrow \(\iota\) in \([1]\) and an arrow \(\beta : a \to b\) in \(B\) such that \(f\beta = ev_{[1], C}(S, \iota)\), then there exists a unique square \(R\) in \(B\) such that \(fR = S\) and
ev_{[1],B}(B,ι) = β. In more elementary terms this says, in the cases where ι is 1_0, 0 → 1 and 1_1, that f enjoys the respective unique lifting properties depicted in

In which dotted arrows in B are the unique liftings of the corresponding arrows down in C. For instance, in the case ι = 1_0 depicted on the left, this says that given β in B and a square in C whose left vertical edge is fβ, one can lift that square uniquely to a square in B whose left vertical edge is β. Clearly one has this unique lifting property when f is a discrete opfibration, and similarly one has the unique lifting property depicted on the right (ι = 1_1) when f is a discrete fibration.

As for the unique lifting property depicted in the middle (ι = 0 → 1), one has this whenever f is a discrete Conduché fibration, that is, when f satisfies the unique lifting of factorisations. As is well-known and easy to check directly, both discrete fibrations and discrete opfibrations possess this property.

For X ∈ Cat the arrow category X^{[1]} is a basic 2-categorical limit construction, namely it is the cotensor of X with the category [1]. As such the data of the corresponding limit cone is

\[
\begin{array}{ccc}
X^{[1]} & \xrightarrow{d_X} & \mathbb{I} \\
\downarrow{c_X} & & \downarrow{c_X} \\
X & & \\
\end{array}
\]

in which d_X and c_X are the functors which on objects take domains and codomains respectively, and the component of α_X at β ∈ X^{[1]} is β viewed as an arrow of X. By 2-dimensional universality d_X and c_X are the components of 2-natural transformations d and c : (−)^{[1]} → 1_{\text{Cat}}, and the α_X are the components of a modification α : d → c. Moreover d is cartesian with respect to a functor f : B → C iff f is a discrete opfibration, and c is cartesian with respect to f iff f is a discrete fibration.

Let T be a collection with object set I. Recall that the middle map p_T : E_T → B_T of its corresponding polynomial is a discrete fibration and a discrete opfibration. Since (−)^{[1]} as a right adjoint preserves discrete fibrations and discrete opfibrations, p_T^{[1]} is such and hence exponentiable. Thus applying (−)^{[1]} componentwise to T’s corresponding polynomial and identifying I = I^{[1]}, gives another polynomial as on
the left
(9)

\[ I \xrightarrow{s_T^1} E_T^1 \xrightarrow{p_T^1} B_T^1 \xrightarrow{t_T^1} I \]

and since \( d \) and \( c \) are cartesian with respect to \( p_T \), one in fact has a 3-cell \((\alpha_B^T, \alpha_E^T)\) in \( \text{Poly}_{\text{Cat}} \) as on the right.

**Definition 5.4.** The full sub-2-category of \( \text{Poly}_{\text{Cat}}(I, I) \) consisting of those polynomials \( I \rightarrow I \) whose middle map is a discrete fibration and a discrete opfibration is denoted as \( \mathcal{D}_I \).

Since discrete fibrations and discrete opfibrations are pullback stable and closed under composition, the monoidal structure \( \text{Poly}_{\text{Cat}}(I, I) \) has by virtue of the composition of polynomials, restricts to \( \mathcal{D}_I \). Applying \( (\cdot)^{[1]} \) componentwise to such polynomials as above is the effect on objects of an endo-2-functor

\[ \mathcal{I}_I : \mathcal{D}_I \rightarrow \mathcal{D}_I. \]

The components of \( d \) and \( c \) give 2-natural transformations \( \mathcal{D}_I \) and \( \mathcal{C}_I : \mathcal{I}_I \rightarrow 1_{\mathcal{D}_I} \), and the components of \( \alpha \) give a modification \( \mathcal{A}_I : \mathcal{D}_I \rightarrow \mathcal{C}_I \). The component of \( \mathcal{A}_I \) with respect to the polynomial corresponding to a collection \( T \) is \((\alpha_B^T, \alpha_E^T)\) depicted in (9).

As is well-known, for a symmetric monoidal closed category \( \mathcal{V} \), the basic notions of monoidal category, lax and strong monoidal functor and monoidal natural transformation admit evident \( \mathcal{V} \)-enriched analogues. In particular the endohoms of \( \text{Poly}_{\text{Cat}} \), or indeed those of any 2-bicategory, and the 2-categories \( \mathcal{D}_I \) defined above are monoidal 2-categories in the straightforward \( \text{Cat} \)-enriched sense. Given monoidal 2-categories \( \mathcal{X} \) and \( \mathcal{Y} \), lax monoidal 2-functors \( F \) and \( G : \mathcal{X} \rightarrow \mathcal{Y} \), and monoidal 2-natural transformations \( \phi \) and \( \psi : F \rightarrow G \), one has an evident notion of *monoidal modification* \( \zeta : \phi \rightarrow \psi \), which is a modification such that

\[
\begin{align*}
\phi \otimes \phi & \rightarrow \phi \\
\psi \otimes \psi & \rightarrow \psi
\end{align*}
\]

commute, where \( \overline{\phi}, \overline{\psi} \) and \( \overline{\phi} \otimes \overline{\psi} \) are the lax monoidal coherences of \( F \) and \( G \).

With these 3-cells monoidal 2-categories form a (strict) 3-category \( \text{Mon-2-CAT} \) equipped with a forgetful 3-functor into \( \text{2-CAT} \). As in the unenriched setting...
monoidal 2-functors $1 \to \mathcal{X}$ may be identified with (strict) monoids in $\mathcal{X}$. We denote by $\text{Mon}(\mathcal{X})$ the 2-category $\text{Mon-2-CAT}(1, \mathcal{X})$, whose objects are monoids and monoid morphisms, and whose 2-cells we refer to as monoid 2-cells. Also of interest for us are weaker notions of monoid morphism, so we denote by $\text{Mon}(\mathcal{X})_c$ and $\text{Mon}(\mathcal{X})_{ps}$ the 2-categories whose objects are monoids in $\mathcal{X}$, and morphisms are lax, colax and pseudo morphisms of monoids whose coherence data is of the form

\[
\begin{array}{ccc}
I & \overset{u_X}{\rightarrow} & X \overset{m_X}{\rightarrow} X \otimes_2 X \\
\downarrow & & \downarrow \cong \\
Y & \overset{f}{\rightarrow} & Y \otimes_2 Y
\end{array}
\quad \begin{array}{ccc}
I & \overset{u_X}{\rightarrow} & X \overset{m_X}{\rightarrow} X \otimes_2 X \\
\downarrow & & \downarrow \cong \\
Y & \overset{f}{\rightarrow} & Y \otimes_2 Y
\end{array}
\quad \begin{array}{ccc}
I & \overset{u_X}{\rightarrow} & X \overset{m_X}{\rightarrow} X \otimes_2 X \\
\downarrow & & \downarrow \cong \\
Y & \overset{f}{\rightarrow} & Y \otimes_2 Y
\end{array}
\]

respectively satisfying the usual axioms, and monoid 2-cells between these. In the case where $\mathcal{X} = \text{Cat}$ with cartesian tensor product, one refinds the usual notions of lax, colax and strong monoidal functor between strict monoidal categories, and monoidal natural transformations between these.

In particular, for any 2-category $\mathcal{K}$ one has the monoidal 2-category $\text{End}(\mathcal{K})$ of endo-2-functors of $\mathcal{K}$. For $\mathcal{X} = \text{End}(\mathcal{K})$ the 2-categories $\text{Mon}(\mathcal{X})$, $\text{Mon}(\mathcal{X})_1$, $\text{Mon}(\mathcal{X})_c$ and $\text{Mon}(\mathcal{X})_{ps}$ just defined are denoted $\text{Mnd}(\mathcal{K})$, $\text{Mnd}(\mathcal{K})_1$, $\text{Mnd}(\mathcal{K})_c$ and $\text{Mnd}(\mathcal{K})_{ps}$ respectively. The 2-cells in any of these 2-categories will be called 2-monad 2-cells.

For any monoid $M$ in a monoidal 2-category $\mathcal{X}$ and monoidal modification $\zeta$ as in

\[
\begin{array}{c}
1 \quad M \quad \mathcal{X} \quad \phi \\
\alpha \quad \zeta \quad \psi \quad \mathcal{Y}
\end{array}
\]

it follows immediately from the 3-category structure of $\text{Mon-2-CAT}$ that $FM$ and $GM$ are monoids in $\mathcal{Y}$, the components $\phi_M$ and $\psi_M$ are monoid morphisms, and $\zeta_M$ is a monoid 2-cell.

**Proposition 5.5.** In the context of definition 5.4, $I_I : \mathcal{D}_I \to \mathcal{D}_I$ is a strong monoidal 2-functor, $D_I$ and $C_I : I_I \to 1_{\mathcal{D}_I}$ are monoidal 2-natural transformations, and $A_I : D_I \to C_I$ is a monoidal modification.

**Proof.** Since $I^{[1]} = I$ one has $I_I(1_I) = 1_I$. Given $P$ and $Q \in \mathcal{D}_I$, since $(-)^{[1]}$ preserves pullbacks and distributivity pullbacks by lemma 5.3, one has a coherence isomorphism $I_I(P) \circ I_I(Q) \cong I_I(P \circ Q)$ induced by the universal properties of the pullbacks and distributivity pullbacks involved in the formation of $I_I(P) \circ I_I(Q)$. The monoidal 2-functor coherence axioms follow because of the uniqueness inherent in how these coherence isomorphisms were induced. This uniqueness, together with the naturality of $d_X : X^{[1]} \to X$, $c_X : X^{[1]} \to X$ and $\alpha_X : d_X \to c_X$ in $X \in \text{Cat}$, enables one to establish that $D_I$ and $C_I$ are monoidal 2-natural transformations and $A_I$ is a monoidal modification. \qed

For an operad $T$ with object set $I$ this last result ensures that $(\alpha_{BT}, \alpha_{ET})$ depicted in (9) underlies a monoid 2-cell in $\mathcal{D}_I$. Since the effect on homs of $\text{P}_{\text{Cat}}$
gives a strong monoidal 2-functor $\mathcal{D}_I \to \text{End}(\text{Cat}/I)$, $\mathcal{P}_{\text{Cat}}$ carries $(\alpha_{B_T}, \alpha_{E_T})$ to a 2-cell in $\text{Mnd}(\text{Cat}/I)$.

**Definition 5.6.** The 2-monad 2-cell just described is denoted

\[
\begin{array}{ccc}
T^{[1]}_{\Sigma} & \xrightarrow{d_T} & T \times T \\
\downarrow \alpha_T & & \downarrow \alpha_T \\
& & 
\end{array}
\]

We now turn to giving an explicit description of the components of $\alpha_T$. Recall from lemma 3.10 that for $X \in \text{Cat}/I$, an arrow of $TX$ is of the form $(\rho, (\gamma_j)_j) : (\alpha \rho, (x_j)_j) \to (\alpha, (y_j)_j)$ where $\alpha : (i_j)_{1 \leq j \leq n} \to i$ is from $T$, $\rho \in \Sigma_n$ and $\gamma_j : x_j \to y_{\rho j}$ is from $X_{\rho j}$. Recall also from the proof of theorem 4.8 that $(\rho, (\gamma_j)_j)$ is levelwise when $\rho$ is an identity, and permutative when the $\gamma_j$ are all identities. We denote by $T^{[1]}_{\Sigma}X$ the subcategory of $(TX)^{[1]}$ consisting of the permutative maps.

Given a subcategory $S$ of a category $C$ which contains all the objects, or equivalently a class of maps of $C$ containing all identities and closed under composition, we denote by $C^S_{\Sigma}$ the full subcategory of the arrow category of $C$ consisting of the $f \in S$. One has functors $d_{c,S}$ and $c_{c,S} : C^S_{\Sigma} \to C$ which on objects take domains and codomains of morphisms of $S$ respectively, and a natural transformation $\alpha_{c,S} : d_{c,S} \to c_{c,S}$ whose component at $f \in C^S_{\Sigma}$ is $f$ viewed as an arrow of $C$.

**Lemma 5.7.** If $T$ is an operad with object set $I$ and $X \in \text{Cat}/I$, then $T^{[1]}_{\Sigma}X = (TX)^{[1]}_{T^I_X}$, $d_{T^I_X} = d_{TX,T^I_X}$, $c_{T^I_X} = c_{TX,T^I_X}$ and $\alpha_{T^I_X} = \alpha_{TX,T^I_X}$.

In particular, $T^{[1]}_{\Sigma}X$ is the full subcategory of the arrow category of $TX$ consisting of the permutative maps.

**Proof.** By the general calculation of example 2.2 applied to the polynomial $(s^{[1]}_T, p^{[1]}_T, t^{[1]}_T)$, an object of $T^{[1]}_{\Sigma}X$ is a pair $(b, h)$ where $b$ is a morphism of $B_T$ and $h : (p^{[1]}_T)^{-1}\{b\} \to X$ is a functor over $I$. Such a morphism $b$ is of the form $\rho : \alpha \rho \to \alpha$, where $\alpha : (i_j)_{1 \leq j \leq n} \to i$ is in $T$ and $\rho \in \Sigma_n$, and in these terms an element of the fibre $(p^{[1]}_T)^{-1}\{b\}$ is a morphism of $E_T$ of the form $\rho : (\alpha \rho, j) \to (\alpha, \rho j)$ where $1 \leq j \leq n$. Thus an object of $T^{[1]}_{\Sigma}X$ consists of $\rho : \alpha \rho \to \alpha$ and $(x_j)_j$ where $x_j \in X_{\rho j}$ for $1 \leq j \leq n$, and one can identify this with the permutative map $(\rho, (1_x)_j)$ in $T^{[1]}_{\Sigma}X$. Similarly unpacking the morphisms of $T^{[1]}_{\Sigma}X$ following the general scheme of example 2.2 one identifies a morphism between permutative maps in $T^{[1]}_{\Sigma}X$ as morphism between them in the arrow category $(TX)^{[1]}$. Having established this explicit description of $T^{[1]}_{\Sigma}$, one obtains those for $d_{T^I_X}$, $c_{T^I_X}$ and $\alpha_{T^I_X}$, using the explicit description of the 2 and 3-cell map of $\mathcal{P}_{\text{Cat}}$ which involves inducing arrows from the universal properties of pullbacks and distributivity pullbacks. \[\square\]

The conical colimit of a diagram

\[
\begin{array}{ccc}
A & \xrightarrow{f} & B \\
\downarrow \phi & & \downarrow g \\
\downarrow \alpha & & 
\end{array}
\]

in a 2-category $K$ is called a **couniversal**. Thus a cocone for (10) consists of $q : B \to Q$ such that $q \phi = \text{id}$, and such a cocone exhibits $Q$ as the couniversal of this diagram when for all $X \in K$, composition with $q$ induces an isomorphism of
categories between $\mathcal{K}(Q, X)$ and the full subcategory of $\mathcal{K}(B, X)$ consisting of those 1-cells $h : B \to X$ such that $h\phi$ is an identity. In other words $q$, is the universal 1-cell which by post-composition makes $\phi$ into an identity. It is more common to consider the coinverter of the above diagram, in which $q$ is the universal 1-cell which by post-composition makes $\phi$ into an isomorphism.

When $\phi$ has a 1-section, that is there is a morphism $i : B \to A$ such that $\phi i = \text{id}$, including or not including the data of $i$ in (10) clearly does not affect the resulting colimit, and in this case $(Q, q)$ is said to be a reflexive coinverter. It is completely straightforward to adapt lemma 2.1 of [13], which considers the case of reflexive coinverters, to exhibit an analogue of the $(3 \times 3)$-lemma for reflexive coinverters. From this it follows that in $\textbf{Cat}$, reflexive coinverters commute with finite products, and so reflexive coinverters are a type of sifted colimit.

**Example 5.8.** For any category $X$ the discrete category of its connected components can be obtained as a coinverter

$$
\begin{array}{ccc}
X & \xrightarrow{d_X} & X^{[1]} \xrightarrow{\pi_0} \pi_0 X.
\end{array}
$$

This coinverter is reflexive via the functor $i_X : X \to X^{[1]}$ given on objects by $i_X x = 1_x$. Thus one recovers the fact that $\pi_0 : \textbf{Cat} \to \textbf{Set}$ preserves finite products from the siftedness of reflexive coinverters.

Let $\mathcal{X}$ be a monoidal 2-category and $\mathcal{S}$ be a set of objects of $\mathcal{X}$. Denote by $\mathcal{S}^*$ the set of tensor products of objects of $\mathcal{S}$, that is, the objects of the smallest monoidal subcategory of $\mathcal{X}$ containing $\mathcal{S}$. Consider the diagram

$$
\begin{array}{ccc}
A & \xrightarrow{f} & B & \xrightarrow{g} & Q
\end{array}
$$

in $\mathcal{X}$.

**Definition 5.9.** We say that $q$ exhibits $Q$ as an $\mathcal{S}$-stable coinverter when for all $X$ and $Y \in \mathcal{S}^*$, the 2-functor

$$
X \otimes (-) \otimes Y : \mathcal{X} \to \mathcal{X}
$$

sends (11) to a coinverter. When $\phi$ has a 1-section then (11) is said to be an $\mathcal{S}$-stable reflexive coinverter.

Taking $X = Y = I$ one finds that for any $\mathcal{S}$, an $\mathcal{S}$-stable coinverter is a coinverter.

**Definition 5.10.** When the 2-cell $\phi$ lives in $\text{Mon}(\mathcal{X})$, then (11) is said to be a monoidally stable coinverter when it is an $\{A, B\}$-stable coinverter in $\mathcal{X}$. When $\phi$ has a 1-section in $\mathcal{X}$, then (11) is said to be a monoidally stable reflexive coinverter.

Note that in the context of definition 5.10, while there are monoid structures on $A$ and $B$, there is not apriori a monoid structure on $Q$. However one has

**Proposition 5.11.** If (11) is a monoidally stable reflexive coinverter in a monoidal 2-category $\mathcal{X}$, then there exists a unique monoid structure on $Q$ making
it a coidentifier in Mon($\mathcal{X}$). Moreover this coidentifier is preserved by any of the inclusions

$$\text{Mon}(\mathcal{X}) \hookrightarrow \text{Mon}(\mathcal{X})_l \quad \text{Mon}(\mathcal{X}) \hookrightarrow \text{Mon}(\mathcal{X})_c \quad \text{Mon}(\mathcal{X}) \hookrightarrow \text{Mon}(\mathcal{X})_{ps}.$$  

\textbf{Proof.} We begin by verifying that

\begin{equation} \tag{12} 
A^\otimes n \xrightarrow{f^\otimes n} B^\otimes n \xrightarrow{g^\otimes n} Q^\otimes n 
\end{equation}

is an $\{A, B\}$-stable coidentifier by induction on $n$. In the base case $n = 0$, (12) is a constant diagram, all constant diagrams of this form are clearly coidentifiers, and so (12) in this case is an absolute coidentifier. The inductive step follows since by the inductive hypothesis, for all $X$ and $Y$ in $\{A, B\}^*$ the first two rows and first two columns of the evident diagram

\begin{align*}
X \otimes A \otimes A^\otimes n \otimes Y & \xrightarrow{=} X \otimes A \otimes B^\otimes n \otimes Y \\
X \otimes B \otimes A^\otimes n \otimes Y & \xrightarrow{=} X \otimes B \otimes B^\otimes n \otimes Y \\
X \otimes Q \otimes A^\otimes n \otimes Y & \xrightarrow{=} X \otimes Q \otimes B^\otimes n \otimes Y \\
X \otimes Q \otimes Q^\otimes n \otimes Y & \xrightarrow{=} X \otimes Q \otimes Q^\otimes n \otimes Y
\end{align*}

are reflexive coidentifiers, and the $(3 \times 3)$-lemma. Thus the unit and multiplication of $Q$ is induced from those of $A$ and $B$ in the evident manner

\begin{align*}
\begin{array}{c}
\xymatrix{ I \ar[r] & I \\
A \ar[r] & B \\
A & B \\}
\end{array}
\quad \begin{array}{c}
\xymatrix{ A \otimes A \ar[r] & B \otimes B \\
A \ar[r] & B \\
A & B \\}
\end{array}
\end{align*}

and the monoid axioms for $Q$ follow easily from the 1-dimensional universal properties of (12).

Suppose that $C$ is a monoid in $\mathcal{X}$, and $h : B \to C$ is a lax morphism of monoids such that $h\phi = \text{id}$ in Mon($\mathcal{X}$)$_{ps}$. By the way that composition works in Mon($\mathcal{X}$)$_{ps}$, this last equation amounts to the equation $h\phi = \text{id}$ in $\mathcal{X}$. By the 1-dimensional universal property of $q$, one has $k : Q \to C$ unique such that $kq = h$. The lax morphism coherence 2-cells $(\overline{k_0}, \overline{k_2})$ for $k$ are induced from the corresponding coherences for $h$ via the 2-dimensional universal properties of $q^\otimes 0 = 1_f$ and $q^\otimes 2$ as in

\begin{align*}
\begin{array}{c}
\xymatrix{ I \ar[r] & I \\
A \ar[r] & B \\
A & B \\}
\end{array}
\quad \begin{array}{c}
\xymatrix{ A^\otimes 2 \ar[r] & B^\otimes 2 \\
A \ar[r] & B \\
A & B \\}
\end{array}
\end{align*}
and the lax morphism coherence axioms for \((\mathcal{K}_0, \mathcal{K}_2)\) follow from those for \((\mathcal{H}_0, \mathcal{H}_2)\) and the 2-dimensional universal properties of \((\mathcal{K})\). In this way \((\mathcal{H})\) has the 1-dimensional universal property of a coidentifier in \(\text{Mon}(\mathcal{X})\), and reversing the 2-cells in this discussion exhibits the 1-dimensional universal property of a coidentifier in \(\text{Mon}(\mathcal{X})\). By the 2-dimensional universal properties of \((\mathcal{H})\) again, \(\mathcal{H}_0\) and \(\mathcal{H}_2\) are isomorphisms or identities iff \(\mathcal{K}_0\) and \(\mathcal{K}_2\) are, and so \((\mathcal{H})\) also has the 1-dimensional universal property of a coidentifier in \(\text{Mon}(\mathcal{X})\). In a similar manner, the 2-dimensional universal properties of \((\mathcal{H})\) in the various 2-categories of monoids under consideration are verified directly using the 2-dimensional universal properties of \((\mathcal{H})\).

Example 5.12. For any 2-category \(\mathcal{K}\) and any endo-2-functor \(R\) of \(\mathcal{K}\), the 2-functor on the left
\[
(-) \circ R : \text{End}(\mathcal{K}) \to \text{End}(\mathcal{K})
\]
and the 2-functor on the right preserves any colimits that \(R\) preserves. Thus given 2-monads \(S\) and \(T\) on \(\mathcal{K}\) which preserve reflexive coidentifiers, a 2-monad 2-cell \(\phi\) between them with a 1-section in \(\text{End}(\mathcal{K})\), and a coidentifier
\[
\begin{array}{ccc}
S & \xrightarrow{\phi} & T \\
\downarrow & & \downarrow \\
Q & & Q
\end{array}
\]
of \(\phi\) in \(\text{End}(\mathcal{K})\), then by proposition 5.11 there is a unique 2-monad structure on \(Q\) making \(q\) the coidentifier of \(\phi\) in \(\text{Mnd}(\mathcal{K})\).

Remark 5.13. In the situation of definition 5.6 the 2-monads \(T\) and \(T^{[1]}_\Sigma\) preserve sifted colimits by \([26]\) theorem 4.4.1. Moreover since \(p_T\) is a discrete fibration it reflects identities, and so the square in
\[
\begin{array}{ccc}
I & \xleftarrow{s_T^{[1]}} & E_T^{[1]} \\
& \searrow & \downarrow \leftarrow \\
& & E_T \\
& \downarrow & \downarrow \leftarrow \\
& & B_T \\
& \searrow & \downarrow \leftarrow \\
& & B_T^{[1]} \\
& \downarrow & \downarrow \leftarrow \\
& & I
\end{array}
\]
is easily verified to be a pullback. Thus this diagram exhibits a 1-section of the 2-cell \((\alpha_{B_T}, \alpha_{E_T})\) in \(\text{Poly}_{\text{Cat}}(I, I)\), and so the situation of definition 5.6 conforms to that of example 5.12. Hence the coidentifier of the 2-monad 2-cell \(\alpha_T\) exists, and is computed as in \(\text{End}(\text{Cat}/I)\).

Definition 5.14. The coidentifying monad morphism of remark 5.13 is denoted
\[
g_T : T \to T/\Sigma.
\]

Remark 5.15. When \(X \to I\) in \(\text{Cat}/I\) is discrete all of \(TX\)’s morphisms are permutative, and so by lemma 5.7, \(T^{[1]}_{\Sigma}X = (TX)^{[1]}\), whence \(T/\Sigma(X) = \pi_0(TX)\) by example 5.8. Thus the 2-monad \(T/\Sigma\) just defined restricts to a monad on \(\text{Set}/I\).

In the remainder of this section we establish analogues of theorems 4.8, 4.12 and 4.15, which say that the algebras of \(T/\Sigma\) correspond to the commutative variants of the algebras of \(T\). These results are established by using endomorphism 2-monads and their variants \([9, 12, 17]\), the details of which we shall recall as needed.
Let $\mathcal{K}$ be a complete 2-category, and for $A \in \mathcal{K}$ and $X \in \textbf{Cat}$ we denote by $A^X$ the cotensor of $X$ with $A$, whose universal property gives isomorphisms as on the left

$$\mathcal{K}(B, A^X) \cong \text{Cat}(X, \mathcal{K}(B, A)) \quad \langle A, B \rangle C = B^{\mathcal{K}(C, A)}$$

2-naturally in $B$. Then the formula on the right for $A$, $B$ and $C \in \mathcal{K}$ defines $\langle A, B \rangle \in \text{End}(\mathcal{K})$. The assignation $B \mapsto \langle A, B \rangle$ is the object map of a 2-functor

$$(A, -) : \mathcal{K} \to \text{End}(\mathcal{K}) \quad \varepsilon_{A,B} : \langle A, B \rangle A \to B$$

which is right adjoint to the 2-functor given by evaluating at $A$, and the counit of this adjunction is as denoted on the right in previous display. As explained in [7, 8], $\langle A, B \rangle$ is thus the hom for an enrichment of $\mathcal{K}$ in $\text{End}(\mathcal{K})$. The units $\eta_A : 1_{\mathcal{K}} \to \langle A, A \rangle$ and compositions $c_{ABC} : \langle B, C \rangle \langle A, B \rangle \to \langle A, C \rangle$ for this enrichment are the unique 2-natural transformations making

\[
\begin{array}{ccc}
A & \xrightarrow{\eta_A} & \langle A, A \rangle \\
\downarrow {\varepsilon_{A,A}} & & \downarrow {\varepsilon_{A,A}} \\
A & \rightarrow & \langle A, A \rangle
\end{array}
\quad
\begin{array}{ccc}
\langle B, C \rangle \langle A, B \rangle & \xrightarrow{c_{ABC,A}} & \langle A, C \rangle A \\
\downarrow {\varepsilon_{B,C}} & & \downarrow {\varepsilon_{B,C}} \\
\langle B, C \rangle B & \rightarrow & C
\end{array}
\]

commute. In particular for $A \in \mathcal{K}$, one has the corresponding endomorphism 2-monad $\langle A, A \rangle$.

By the universal property of $\varepsilon_{A,A}$ one has a bijection between 2-natural transformations $\phi : T \to \langle A, A \rangle$ and morphisms $a : TA \to A$ in $\mathcal{K}$, and when $T$ is a 2-monad, $\phi$ satisfies the axioms of a strict morphism of 2-monads iff $a$ satisfies the axioms of a strict $T$-algebra. As Kelly first observed in [9], this correspondence extends to give a description of lax and pseudo algebra structures in terms of morphisms of 2-monads. In the lax case the data for a lax morphism of 2-monads $T \to \langle A, A \rangle$ includes an underlying 2-natural transformation $\phi$ as above, together with modifications $\phi_0$ and $\phi_2$ as in

\[
\begin{array}{ccc}
1_{\mathcal{K}} & \xrightarrow{\eta} & T \\
\downarrow {\phi} & & \downarrow {\phi} \\
\langle A, A \rangle & \xrightarrow{T^2} & \langle A, A \rangle^2
\end{array}
\quad
\begin{array}{ccc}
A & \xrightarrow{\eta_A} & TA \\
\downarrow {\phi} & & \downarrow {\phi} \\
A & \xrightarrow{T^2A} & TA
\end{array}
\]

which by the 2-dimensional part of the universal property of $\varepsilon_{A,A}$, corresponds to 2-cells $\pi_0$ and $\pi_2$ in $\mathcal{K}$. Moreover the lax morphism coherence axioms on $\phi_0$ and $\phi_2$ correspond to the lax algebra coherence axioms on $\pi_0$ and $\pi_2$, and $\phi_0$ and $\phi_2$ are invertible iff $\pi_0$ and $\pi_2$ are. In this way, lax or pseudo morphisms $T \to \langle A, A \rangle$ of 2-monads, may be identified with lax or pseudo $T$-algebra structures on $A$.

**Theorem 5.16.** Let $T$ be an operad with object set $I$ and let $H \to I$ be an object of $\textbf{Cat}/I$. To give $H$ the structure of a lax, pseudo or strict $(T/\Sigma)$-algebra is to give it the structure of a commutative lax, commutative pseudo or commutative strict morphism $H : T \to \textbf{Cat}$, respectively.

**Proof.** By definition 5.14 and proposition 5.11 to give a lax morphism $T/\Sigma \to \langle H, H \rangle$ of 2-monads is to give a lax morphism $T \to \langle H, H \rangle$ such that underlying 2-natural transformation $\phi$ post-composes with $\alpha_T$ of definition 5.6 to give an identity.
By the universal property of \( \varepsilon_{H,H} \) this is the same as giving a lax \( T \)-algebra structure on \( H \), whose 1-cell datum \( a : TH \to H \) post-composes with \( \alpha_{T,H} \) to give an identity. By theorem 4.8 a lax \( T \)-algebra structure on \( H \) is the same thing as a lax morphism \( H : T \to \text{Cat} \), and by the explicit description of \( \alpha_{T,H} \) provided by lemma 5.7, the condition that \( a : TH \to H \) post-composes with \( \alpha_{T,H} \) to give an identity corresponds to the condition that the symmetries of \( H : T \to \text{Cat} \) are identities. The pseudo and strict cases follow in the same way by considering pseudo and strict morphisms of 2-monads \( T \to \langle H, H \rangle \).

As Kelly also understood [9], the different types of morphisms of algebras of a 2-monad can also be regarded as morphisms of 2-monads. In the above setting of a complete 2-category \( K \), given \( f : A \to B \) in \( K \) and \( T \in \text{End}(K) \), data \( \langle \hat{a}, \hat{b}, \hat{\phi} \rangle \) in \( \text{End}(K) \) is in bijection with the data \( \langle a, b, \phi \rangle \) in \( K \) by the universal property of \( \varepsilon_{A,B} \). Let \( \{ f, f \} \) be the comma object \( \langle A, f \rangle \downarrow \langle f, B \rangle \), denote its defining comma square in \( \text{End}(K) \) as on the left

\[
\begin{array}{ccc}
T & \xrightarrow{\hat{b}_f} & \langle B, B \rangle \\
\hat{a}_f & \Downarrow & \phi \\
\langle A, A \rangle & \xrightarrow{\langle f, f \rangle} & \langle A, B \rangle \\
\end{array}
\]

and denote the corresponding data in \( K \) as on the right. In terms of this data in \( K \), the 1-dimensional part of the comma object universal property says that given \( \langle a, b, \phi \rangle \) as above, one has a unique 2-natural transformation \( \phi' : T \to \{ f, f \} \) such that \( a_f \phi'_A = a, b_f \phi'_B = b \) and \( \phi_f \phi'_A = \phi \). In particular from

\[
\begin{array}{ccc}
A & \xrightarrow{1_A} & A \\
\downarrow & \downarrow & \downarrow \\
B & \xrightarrow{1_B} & B \\
\end{array}
\]

and

\[
\begin{array}{ccc}
\{ f, f \} & \xrightarrow{a_f} & A \\
\phi_f & \Downarrow & \phi_f \\
\{ f, f \} & \xrightarrow{b_f} & B \\
\end{array}
\]

one induces \( \eta^f : 1_K \to \{ f, f \} \) and \( \mu^f : \{ f, f \} \to \{ f, f \} \) unique providing the unit and multiplication of a 2-monad, making \( a_f \) and \( b_f \) into strict algebra structures for this 2-monad, and \( \{ f, \phi_f \} \) into a lax morphism between them. Moreover, with respect to this 2-monad structure on \( \{ f, f \} \) and the endomorphism 2-monad structures on \( \langle A, A \rangle \) and \( \langle B, B \rangle \), \( \hat{a}_f \) and \( \hat{b}_f \) become strict morphisms of 2-monads.

From the universal property of \( \varepsilon_{A,B} \) it follows that to give a lax, pseudo or strict morphism \( T \to \{ f, f \} \) of 2-monads, is the same as giving \( A \) and \( B \) lax, pseudo or strict algebra structures respectively, and a 2-cell \( \phi \) providing the coherence making \( \{ f, \phi \} \) a lax morphism of \( T \)-algebras. Composing with \( \hat{a}_f \) and \( \hat{b}_f \) one recovers the \( T \)-algebra structures on \( A \) and \( B \) as morphisms of 2-monads. To
summarise, \( \{f,f\}_1 \) is a 2-monad on \( K \) which classifies lax morphisms of algebras with underlying 1-cell \( f \). Replacing the comma object \( \langle A, f \rangle \downarrow \langle f, B \rangle \) in the above discussion by either \( \langle f, B \rangle \downarrow \langle A, f \rangle \), the isocomma object or the pullback, produces the 2-monads \( \{f,f\}_c \), \( \{f,f\}_p \) and \( \{f,f\} \), which similarly classify colax, pseudo and strict morphisms of algebras with underlying 1-cell \( f \) respectively.

\textbf{Theorem 5.17.} Let \( T \) be an operad with object set \( I \) and let \( f : H \to K \) be a morphism of \( \text{Cat}/I \). Suppose also that one has the structure of commutative lax morphism \( T \to \textbf{Cat} \) on both \( H \) and \( K \). Then to give \( f \) the structure of lax, colax, pseudo or strict \( (T/\Sigma)\)-algebra morphism, is to give \( f \) the structure of lax-natural, colax-natural, pseudo-natural or natural transformation respectively.

\textbf{Proof.} To give \( f \) the structure of a lax \( (T/\Sigma)\)-algebra morphism is to give a lax morphism \( (T/\Sigma) \to \{f,f\}_1 \) of 2-monads whose composites with \( \tilde{a}_f \) and \( \tilde{b}_f \) correspond as monad morphisms to the given \( T/\Sigma \)-algebra structures on \( H \) and \( K \). By definition \ref{5.14} and proposition \ref{5.11}, to give a lax morphism \( (T/\Sigma) \to \{f,f\}_1 \), is to give a lax morphism \( \phi : T \to \{f,f\}_1 \) whose underlying 2-natural transformation post composes with \( \alpha_T \) to an identity. This last condition is equivalent, by the 2-dimensional universal property of the defining comma square for \( \{f,f\}_1 \) in \( \text{End}(K) \), to the condition that the underlying 2-natural transformations of \( \alpha_f \phi \) and \( \tilde{b}_f \) post composes with \( \alpha_T \) to identities, but this just says in turn that the underlying \( T \)-algebra structures on \( H \) and \( K \) correspond to commutative lax morphisms \( T \to \textbf{Cat} \). Thus the result follows from theorem \ref{4.12} in the lax case. For colax, pseudo and strict \( T/\Sigma \)-algebra morphisms, one argues in the same way using \( \{f,f\}_c \), \( \{f,f\}_p \) and \( \{f,f\} \) respectively. \( \square \)

One also has 2-monads that classify algebra 2-cells. In the situation of a complete 2-category \( K \), 1-cells \( f \) and \( g : A \to B \), and a 2-cell \( \gamma : f \to g \), one can define the comma object

\[
\begin{CD}
\{\gamma,\gamma\}_1 @>b_f>> \langle B, B \rangle \\
\langle A, A \rangle @<a_f<< \langle A, B \rangle^{[1]} \Downarrow \langle \phi, B \rangle
\end{CD}
\]

in \( \text{End}(K) \). In a similar manner to our discussion of \( \{f,f\}_1 \) above, one can then exhibit the unit and multiplication for a 2-monad structure on \( \{\gamma,\gamma\}_1 \), and describe a bijection between lax morphisms of 2-monads \( T \to \{\gamma,\gamma\}_1 \) and algebra 2-cells, between lax morphisms of lax \( T \)-algebras. As before one classifies such algebra 2-cells between stricter types of algebra by using the corresponding stricter type of morphism of 2-monads \( T \to \{\gamma,\gamma\}_1 \), and one classifies algebra 2-cells between colax, pseudo and strict morphisms by considering the appropriate 2-monad \( \{\gamma,\gamma\}_c \), \( \{\gamma,\gamma\}_p \) and \( \{\gamma,\gamma\} \), obtained by reversing the direction of the comma object, taking an isocomma object or a pullback respectively. The proof of

\textbf{Theorem 5.18.} Let \( T \) be an operad with object set \( I \). Suppose that one has commutative lax morphisms of operads \( H \) and \( K : T \to \textbf{Cat} \), and lax-natural (resp. colax-natural) transformations \( (f,\overline{f}) \) and \( (g,\overline{g}) : H \to K \). Then to give a modification \( (f,\overline{f}) \to (g,\overline{g}) \) is to give an algebra 2-cell between the corresponding lax (resp. colax) morphisms of \( (T/\Sigma) \)-algebras.
then unfolds analogously to that of theorem 5.17.

For each type of 2-category of algebra for 2-monads, a 2-monad morphism $T \to S$ induces a “forgetful” 2-functor from the corresponding 2-category of algebras of $S$ to those of $T$. In the present situation of $q_T : T \to T/\Sigma$, theorems 5.16-5.18 say that these correspond to the inclusions of the commutative lax, pseudo or strict operad morphisms $T \to \textbf{Cat}$ amongst the general such morphisms. Regardless of which type of algebra 1-cells are considered, these inclusions are clearly 2-fully faithful.

At the beginning of this article we introduced the notation $T/\Sigma$ to denote the monad on $\text{Set}/I$ arising from the operad $T$ via the standard construction of a monad from an operad. In this standard view the effect of $T/\Sigma$ on $X \to I$ in $\text{Set}/I$ is given by the formula

$$ (T/\Sigma(X))_i = \prod_{n \in \mathbb{N}} \left( \prod_{i_1, \ldots, i_n} T((i_j)_j; i) \times \prod_{j=1}^n X_{i_j} \right) / \Sigma_n $$

interpreted as follows. The term in the bracket is acted on by $\Sigma_n$ by permuting the variables $(i_1, \ldots, i_n)$, and then the notation $(-)/\Sigma_n$ is the standard notation for identifying the orbits of this $\Sigma_n$-action. By remark 5.15 $T/\Sigma$ given in definition 5.14 restricts to a monad on $\text{Set}/I$, and by theorems 5.16 and 5.17 its algebras coincide with the version of $T/\Sigma$ defined in the standard way via (13). In this way these two uses of the notation $T/\Sigma$ are consistent. Thus one may say that the defining coidentifier of $T/\Sigma$ of definition 5.14 is a nice conceptual way of expressing the formula (13).

### 6. $\Sigma$-free operads

Let $T$ be a collection with object set $I$. The action of permutations on operations of $T$ provide, for any sequence $(i_j)_{1 \leq j \leq n}$ of objects of $T$, $i \in I$, and any permutation $\rho \in \Sigma_n$, a bijection

$$ (-)\rho : T((i_{\rho j})_j; i) \to T((i_j)_j; i). $$

It can happen that the sequences $(i_{\rho j})_j$ and $(i_j)_j$ are in fact equal, for instance when all the $i_j$’s are the same element of $I$. In such cases one can then ask whether $(-)\rho$ has any fixed points. A collection $T$ is said to be $\Sigma$-free when for all $(i_j)_{1 \leq j \leq n}$, $i$ and $\rho$ as above such that $\rho \neq 1_n$ and $(i_{\rho j})_j = (i_j)_j$, the bijection (14) has no fixed points. A $\Sigma$-free operad is an operad whose underlying collection is $\Sigma$-free.

We now characterise the $\Sigma$-freeness of a collection in various ways. Preliminary to this, it is useful to have various alternative characterisations of those categories equivalent to discrete categories. For any category $X$ we denote by $q_X : X \to \pi_0 X$ the surjective-on-objects functor which sends $x \in X$ to its connected component. In other words, $q_X$ is the coidentifier of the 2-cell which arises from taking the cotensor of $X$ with $[1]$ (as in example 5.8), and moreover is the component at $X$ of the unit of the adjunction with left adjoint $\pi_0 : \textbf{Cat} \to \textbf{Set}$. Recall that a category is indiscrete when it is equivalent to the terminal category 1, or equivalently when it is non-empty and there is a unique morphism between any two objects. The straight forward proof of

**Lemma 6.1.** For $X \in \textbf{Cat}$ the following statements are equivalent.

1. $X$ is equivalent to a discrete category.
(2) $q_X : X \to \pi_0 X$ is fully faithful.
(3) $X$ is a groupoid and every morphism of $X$ is unique in its hom-set.
(4) $X$ is a coproduct of indiscrete categories.

is left to the reader. Categories equivalent to discrete categories are closed under various 2-categorical constructions relevant for us.

Lemma 6.2.  
(1) If $p : E \to B$ is a discrete fibration and $B$ is equivalent to a discrete category, then so is $E$.

(2) In a pullback as on the left

\[
P \ar{r}{f} \ar{d}{g} & B \ar{d}{p} \\
A \ar{r}{f} & C
\]

in $\text{Cat}$, if $A$, $B$ and $C$ are equivalent to discrete categories, then so is $P$.

(3) In a distributivity pullback around $(f, g)$ in $\text{Cat}$ as on the right in the previous display in which $f$ is a discrete fibration, if $A$, $B$ and $C$ are equivalent to discrete categories, then so are $P$ and $Q$.

(4) If in

\[
A \ar{r}{f} \ar{d}[g] & B \ar{d}{q} & Q
\]

$q$ is a reflexive coequalizer of $\phi$ and $B$ is equivalent to a discrete category, then so is $Q$, and $q$ is surjective on objects and arrows.

Proof. Using lemma 6.1(3), together with lemma 2.1 in the case of (3), one easily verifies (1)-(3) directly. In the case of (4) by lemma 6.1(4) the coequalizer diagram decomposes as

\[
\coprod_i A_i \ar{r}{LL_i f_i} & \coprod_i B_i \ar{r}{LL_i q_i} & \coprod_i Q_i
\]

in which each summand is a reflexive coequalizer diagram, and $B_i$ is indiscrete. Thus by lemma 6.1(4) it suffices to consider the case where $B$ is indiscrete. Recall that the functor $(-)_0 : \text{Cat} \to \text{Set}$ which sends every category to its set of objects, has a right adjoint section $\text{ch} : \text{Set} \to \text{Cat}$, which sends every set $X$ to the category whose objects are the elements of $X$, and where there is a unique arrow between any two objects. Given $x_1$ and $x_2 \in X$, we denote the unique arrow $x_1 \to x_2$ in $\text{ch}(X)$ simply as $(x_1, x_2)$. When $B$ is indiscrete one may regard it as $\text{ch}(X)$ for some non-empty set $X$. It suffices to show that in this case the coequalizer is computed as on the left

\[
A \ar{r}{f} \ar{d}[g] & \text{ch}(X) \ar{d}[\text{ch}(q)] & \text{ch}(Q) \\
A_0 \ar{r}[f_0] & X \ar{r}[g_0] & Q
\]

where the diagram on the right is a coequaliser in $\text{Set}$. Let $h : \text{ch}(X) \to C$ be such that $h \phi = \text{id}$. Since $\text{ch}(q)$ is clearly surjective on objects and arrows and thus an epimorphism in $\text{Cat}$, it suffices to show that there exists $h' : \text{ch}(Q) \to C$ such that
$h = h'ch(q)$. The object map $h'_0$ is unique such that $h_0 = h'_0q$ by the coequaliser in Set. Given $y_1$ and $y_2 \in Q$ we must give $h'(y_1, y_2) : h'y_1 \to h'y_2$ in $C$, and this is done by choosing $x_1$ and $x_2$ in $X$ such that $qx_1 = y_1$ and $qx_2 = y_2$, and then defining $h'(y_1, y_2) = h(x_1, x_2)$. The functoriality of $h'$ is immediate from that of $h$ as long as $h'$’s arrow map is well-defined.

To establish this well-definedness we must show that if $qx_1 = qx'_1$ and $qx_2 = qx'_2$, then $h(x_1, x_2) = h(x'_1, x'_2)$. Since $f_0$ and $g_0$ have a common section given by the object map of the 1-section of $\phi$, the equivalence relation on $X \times X$ defined by $(x_1, x_2) \sim (x'_1, x'_2)$ if $qx_1 = qx'_1$ and $qx_2 = qx'_2$ is the smallest equivalence relation which contains

$$\{((x_1, x_2), (x'_1, x'_2)) : \exists a, fa = x_1 \text{ and } ga = x'_1\}$$

$$\cup$$

$$\{((x_1, x_2), (x_1, x'_2)) : \exists a, fa = x_2 \text{ and } ga = x'_2\}.$$  

Thus it suffices to show

- If $x_1, x'_1$ and $x_2 \in X$ such that $\exists a, fa = x_1$ and $ga = x'_1$, then $h(x_1, x_2) = h(x'_1, x_2)$.
- If $x_1, x_2$ and $x'_2 \in X$ such that $\exists a, fa = x_2$ and $ga = x'_2$, then $h(x_1, x_2) = h(x_1, x'_2)$.

In the first of these situations note that one has a triangle as on the left

![Diagram](attachment:image.png)

in $\text{ch}(X)$ sent by $h$ to the desired equality since $h\phi_0 = \text{id}$, and similarly one applies $h$ to the triangle on the right for the other situation.

Recall from section 3 that

\[ I \xrightarrow{\delta_T} E_T \xrightarrow{p_T} B_T \xrightarrow{t_T} I \]

denotes the polynomial corresponding to a collection $T$ with object set $I$.

**Proposition 6.3.** For a collection $T$ with object set $I$ the following statements are equivalent:

1. $T$ is $\Sigma$-free.
2. $B_T$ is equivalent to a discrete category.
3. The naturality square

\[
\begin{array}{ccc}
E_T & \xrightarrow{p_T} & B_T \\
\downarrow q_{E_T} & & \downarrow q_{B_T} \\
\pi_0E_T & \xrightarrow{\pi_0p_T} & \pi_0B_T
\end{array}
\]

is a pullback.

**Proof.** (1)⇒(2): By lemma 6.1 it suffices to show that when $T$ is $\Sigma$-free, that every morphism of $B_T$ is unique in its hom-set. Recall that a morphism of $B_T$ is of the form $\rho : \alpha \rho \to \alpha$, where $\alpha : (i_j)_{1 \leq j \leq n} \to i$ is in $T$ and $\rho \in \Sigma_n$. Thus to give a pair of morphisms with the same domain and codomain, is to give $\alpha$ as above, $\rho_1$
and $\rho_2 \in \Sigma_n$ such that $\alpha \rho_1 = \alpha \rho_2$. Thus $\alpha = \alpha (\rho_2 \rho_1^{-1})$, and so $\Sigma$-freeness implies $\rho_2 \rho_1^{-1} = 1_n$, whence $\rho_1 = \rho_2$.

(2)$\Rightarrow$(3): By lemma 6.2 $E_T$ is also equivalent to a discrete category, and so by lemma 6.1 $q_{E_T}$ and $q_{B_T}$ are surjective-on-objects equivalences. We check that the naturality square is a pullback on objects. Let $b \in B_T$ and $c \in \pi_0 E_T$ such that $q_{B_T} b = \pi_0 p_T c$. Choose $e \in E_T$ such that $q_{E_T} e = c$. Since $p_T e$ and $b$ are in the same connected component, there is a unique isomorphism $p_T e \cong b$, and one has a unique lifting of this to $e \cong e'$ in $E_T$. Thus $e'$ is an object of $E_T$ such that $p_T e' = b$ and $q_{E_T} e' = c$. To see that it is unique, suppose that one has $e_1$ and $e_2$ in $E_T$ such that $q_{E_T} e_1 = q_{E_T} e_2$ and $p_T e_1 = p_T e_2$. By the first of these equations one has a unique isomorphism $e_1 \cong e_2$, and this is sent to an identity by $p_T$. As a discrete fibration $p_T$ reflects identities, and so the isomorphism $e_1 \cong e_2$ is an identity. To say that the naturality square is a pullback on arrows, is to say that given $\beta : b_1 \cong b_2$ in $B_T$ and $c \in \pi_0 E_T$ such that $\pi_0 p_T 1_c = q_{B_T} \beta$, then there is a unique isomorphism $\varepsilon : e_1 \cong e_2$ in $E_T$ such that $q_T \varepsilon = 1$ and $p_T \varepsilon = \beta$. But $e_1$ and $e_2$ are determined uniquely since the square is a pullback on objects, and $\varepsilon$ is determined uniquely since $E_T$ and $B_T$ are equivalent to discrete categories.

(3)$\Rightarrow$(1): We prove the contrapositive. Suppose $T$ is not $\Sigma$-free. Then one has $\alpha : \alpha 1 \leq i \leq \alpha n \Rightarrow i$ in $T$ and $1_n \not\in \Sigma_n$ such that $\alpha \rho = \alpha$. Choose $1 \leq k < l \leq n$ such that $\rho k = l$, for instance by letting $k$ be the least such that $\rho k \neq k$. Then $(\alpha, k)$ and $(\alpha, l)$ are distinct objects of $E_T$, which are in the same connected component since one has $\rho : (\alpha, k) \rightarrow (\alpha, l)$, and one has $p_T (\alpha, k) = \alpha = p_T (\alpha, l)$, and so $q_{E_T}$ and $p_T$ are not jointly monic. 

Thanks to this last result, for a $\Sigma$-free collection $T$ one has the 1-cell $(q_{B_T}, q_{E_T})$

\[
\begin{array}{ccc}
E_T^{[1]} & \xrightarrow{p_T^{[1]}} & B_T^{[1]} \\
\pi_0 E_T & \equiv & \pi_0 B_T \\
I & \xrightarrow{s_T^{[1]}} & E_T & \xleftarrow{t_T^{[1]}} & I \\
\pi_0 s_T & \equiv & \pi_0 s_T & \equiv & \pi_0 t_T \\
\pi_0 E_T & \xrightarrow{q_{E_T}} & B_T & \xleftarrow{q_{B_T}} & \pi_0 E_T \\
\end{array}
\]

(15)

in $\text{Poly}_{\text{Cat}}(I, I)$ whose composite with $(\alpha_{B_T}, \alpha_{E_T})$ is an identity. Composition of polynomials makes $\text{Poly}_{\text{Cat}}(I, I)$ a monoidal 2-category, and the main technical result of this section is

**Lemma 6.4.** $(q_{B_T}, q_{E_T})$ is the coidentifier of $(\alpha_{B_T}, \alpha_{E_T})$ in $\text{Poly}_{\text{Cat}}(I, I)$. This coidentifier is monoidally stable and preserved by $(P_{\text{Cat}})_{I, I} : \text{Poly}_{\text{Cat}}(I, I) \rightarrow \text{End}(\text{Cat}/I)$.

whose proof we defer until after considering its main consequence, which is

**Theorem 6.5.** Let $T$ be a $\Sigma$-free operad with object set $I$. Then $T/\Sigma$ is a polynomial 2-monad whose underlying polynomial is

$$ I \xrightarrow{\pi_0 s_T} \pi_0 E_T \xrightarrow{\pi_0 p_T} \pi_0 B_T \xrightarrow{\pi_0 t_T} I. $$
PROOF. Since $P_{\text{Cat}/I}$ preserves the coidentifier (15), it may be regarded as sending (15) to the defining coidentifier in $\text{End}(\text{Cat})$ of $T/\Sigma$. Since (15) is monoidally stable, by proposition 5.11, the polynomial $(\pi_0 s_T, \pi_0 p_T, \pi_0 t_T)$ acquires a unique 2-monad structure in $\text{Poly}_{\text{Cat}}$ making $(q_{B_T}, q_{E_T})$ a morphism of monads. Since $P_{\text{Cat}}$ is a homomorphism of 2-bicategories, it sends this 2-monad in $\text{Poly}_{\text{Cat}}$ to a 2-monad on $\text{Cat}/I$, whose underlying endo-2-functor coincides with that of $T/\Sigma$ and which makes $q_T : T \to T/\Sigma$ a morphism of 2-monads. However by definition 5.14 and example 5.12, $T/\Sigma$’s 2-monad structure is unique with this last property. □

The remainder of this section is concerned with the proof of lemma 6.4. To this end we denote by $\mathcal{E}_I$ the full sub-2-category of $\text{Poly}_{\text{Cat}}(I,I)$ consisting of those polynomials

$$I \xrightarrow{s} E \xrightarrow{p} B \xrightarrow{t} I$$

such that $p$ is a discrete fibration with finite fibres, and $B$ is equivalent to a discrete category. By lemma 6.2 $E$ is also equivalent to a discrete category.

**Lemma 6.6.** A composite of polynomials in $\mathcal{E}_I$ is in $\mathcal{E}_I$.

**Proof.** We suppose that $(s_1, p_1, t_1)$ and $(s_2, p_2, t_2)$ are in $\mathcal{E}_I$ and form their composite $(s_3, p_3, t_3)$ as in

$$
\begin{array}{ccc}
E_3 & \xrightarrow{p_3} & B_3 \\
| & & | \\
\downarrow{s_3} & \downarrow{pb} & \downarrow{t_3} \\
I & \xrightarrow{p_1} & B_1 \xleftarrow{t_1} I \\
| & & | \\
\downarrow{pb} & \downarrow{dpb} & \downarrow{pb} \\
E_1 & \xrightarrow{p_2} & B_2 \xleftarrow{t_2} I.
\end{array}
$$

The functors which are discrete fibrations, discrete opfibrations and have finite fibres are closed under composition and stable by pullback along arbitrary functors. Thus $p_3$ is such a functor since $p_1$ and $p_2$ are. By lemma 6.2 one deduces successively that $P$, $B_3$, $F$ and $E_3$ are all equivalent to discrete categories. □

We consider now a 2-cell $(\alpha_1, \alpha_2)$ which has a 1-section, and a 1-cell $(q_1, q_2)$

$$
\begin{array}{ccc}
E_3 & \xrightarrow{p_3} & B_3 \\
| & & | \\
\downarrow{pb} & \downarrow{dpb} & \downarrow{pb} \\
E_1 & \xrightarrow{p_1} & B_1 \xleftarrow{p_1} I \\
| & & | \\
\downarrow{q_2} & \downarrow{q_1} & \downarrow{q_2} \\
I & \xrightarrow{p_2} & B_2 \xleftarrow{p_2} I.
\end{array}
$$

(16)

in $\mathcal{E}_I$ such that $(q_1, q_2)(\alpha_1, \alpha_2) = \text{id}.$
Lemma 6.7. If in (16) $q_1$ is the coidentifier in $\text{Cat}$ of $\alpha_1$, then $(q_1,q_2)$ is a coidentifier of $(\alpha_1,\alpha_2)$ in $\text{Poly}_{\text{Cat}}(I,I)$, which is preserved by composition on either side by polynomials from $\mathcal{E}_I$.

Proof. Since $\Sigma_{B_1} : \text{Cat}/B_1 \to \text{Cat}$ creates colimits, one can regard $q$ and $\alpha_1$ as a coidentifier in $\text{Cat}/B_1$. Then $r$ and $\alpha_2$ can be regarded as obtained from these by applying $\Delta_p$, and thus since $p$ is exponentiable, $r$ is the coidentifier of $\alpha_2$ in $\text{Cat}/E_1$, and thus also in $\text{Cat}$ since $\Sigma_{E_1}$ creates colimits.

To show that (16) is a coidentifier in $\text{Poly}_{\text{Cat}}(I,I)$ it suffices to show that given $r_1$, $r_2$ and $f$ as in

making the square $(r_2,f,r_1,p_2)$ a pullback, then the induced square $(s_2,f,s_1,p_1)$ is also a pullback. To check that the induced square is a pullback on objects, consider $b \in B_1$ and $x \in X$ such that $s_1b = fx$. Since $q_1$ is surjective on objects, one has $b_2$ such that $q_1b_2 = b$, and so there is $e_2 \in E_2$ unique such that $p_2e_2 = b_2$ and $r_2e_2 = x$. Thus $e = q_2e_2$ satisfies $p_1e = b$ and $s_2e = x$. To see that $e$ is unique with these properties, consider $e' \in E_1$ such that $p_1e' = b$ and $s_2e' = x$. By the pullback $(q_2,p_1,q_1,p_2)$ there is $e_2' \in E_2$ unique such that $p_2e_2' = b_2$ and $q_2e_2' = e'$. But then $r_2e_2 = r_2e_2'$ and $p_2e_2 = p_2e_2'$, and so since $(p_2,r_2)$ are jointly monic $e_2 = e_2'$, and so $e = e'$. To check that the induced square is a pullback on arrows, given arrows $\beta$ in $B_1$ and $\alpha$ in $X$, one establishes the existence of an arrow $\gamma$ in $E_1$ such that $p_1\gamma = \beta$ and $s_2\gamma = \alpha$ as in the objects case, and the uniqueness follows since $E_1$ is equivalent to a discrete category.

It remains to show that taking a situation (16) in which $q_1$ is a coidentifier of $\alpha_1$ in $\text{Cat}$, and either pre or post-composing it with a polynomial from $\mathcal{E}_I$, gives another such situation. Let $P \in \mathcal{E}_I$ and let the diagram on the left

denote the result of post composing (16) by $P$, and let the diagram on the right denote the result of pre composing (16) by $P$. Our task is to show that in $\text{Cat}$, $q'_1$ is a coidentifier of $\alpha'_1$ and $q''_1$ is a coidentifier of $\alpha''_1$. In the case of $q'_1$ note that $\alpha'_1 = T\alpha_1$ and $q'_1 = Tq_1$ where $T = P_{\text{Cat}I,I}(P)$. By [26] theorem 4.4.1, $T$ preserves sifted colimits and so $q'_1$ is indeed the coidentifier of $\alpha'_1$. In the case of $q''_1$ if one
denotes by

\[
\begin{array}{ccc}
Q & \xrightarrow{d} & R \\
\downarrow c & & \downarrow k \\
\end{array}
\]

the diagram in \(\text{End}(\text{Cat}/I)\) which is the result of applying \(P_{\text{Cat} I I}\) to (16), and by \(X \rightarrow I\) the third map comprising the polynomial \(P\), then \(\alpha''_1\) and \(q''_1\) are the result of evaluating (17) at \(X \in \text{Cat}/I\). Note in particular, that when \(X = 1\) this will be \(\alpha_1\) and \(q_1\), with \(q_1\) a coidentifier of \(\alpha_1\) by hypothesis. In the general case note that one has

\[
\begin{array}{ccc}
QX & \xrightarrow{d_X} & RX \\
\downarrow c_X & & \downarrow k_X \\
Q1 & \xrightarrow{e_1} & R1 \\
\end{array}
\]

One can regard \(k_X\) and \(\zeta_X\) as living in \(\text{Cat}/SX\), and since \(\Sigma_{SX} : \text{Cat}/SX \rightarrow \text{Cat}\) creates colimits, \(k_X\) coidentiﬁes \(\zeta_X\) in \(\text{Cat}/SX\) iff it does so in \(\text{Cat}\). Similarly \(k_1\) may be regarded as a coidentiﬁer of \(\zeta_1\) in \(\text{Cat}/S1\). Since the 2-natural transformations \(d, c\) and \(k\) are cartesian, one may regard the top row as being obtained from the bottom by pulling back along \(S(!)\). Since \(S\) is familial and opfamilial by [26] theorem 4.3.12, it preserves fibrations and opfibrations, whence \(S(!)\) is a ﬁbration and opfibration, and so is exponentiable. Thus \(\Delta_{S(!)}\) preserves colimits, and so \(\zeta_X = \alpha''_1\) is indeed coidentiﬁed by \(k_X = q''_1\).

**Proof.** (of lemma 6.4). The situation (15) clearly conforms to that of (16), and by definition \(q_{B^I_T}\) is the coidentiﬁer of \(\alpha_{B^I_T}\) in \(\text{Cat}\). To show that (15) is a monoidally stable coidentiﬁer, one must show that it is a coidentiﬁer in \(\text{PolyCat}(I, I)\) which is preserved by pre or post composition with composites of

\[
I \xrightarrow{s_T} E_T \xrightarrow{p_T} B_T \xrightarrow{t_T} I \quad I \xrightarrow{s_1^{[1]}} E_T^{[1]} \xrightarrow{p_1^{[1]}} B_T^{[1]} \xrightarrow{t_1^{[1]}} I.
\]

Since these are in \(\mathcal{E}_I\) the result follows by lemmas 6.6 and 6.7.

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