CERTIFYING GLOBAL OPTIMALITY OF AC-OPF SOLUTIONS VIA SPARSE POLYNOMIAL OPTIMIZATION

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Abstract. We report the experimental results on certifying 1% global optimality of solutions of AC-OPF instances from PGLiB via the CS-TSSOS hierarchy — a moment-SOS based hierarchy that exploits both correlative and term sparsity, which can provide tighter SDP relaxations than Shor’s relaxation. Our numerical experiments demonstrate that the CS-TSSOS hierarchy scales well with the problem size and is indeed useful in certifying global optimality of solutions for large-scale real world problems, e.g., the AC-OPF problem. In particular, we are able to certify 1% global optimality for a challenging AC-OPF instance with 6515 buses involving 14398 real variables and 63577 constraints.

1. Introduction

Background on polynomial optimization. Polynomial optimization considers optimization problems where both the cost function and constraints are defined by polynomials, which widely arises in numerous fields, such as optimal power flow [11], numerical analysis [21], computer vision [34], deep learning [8], discrete optimization [26], etc. Even though it is usually not difficult to find a locally optimal solution via a local solver (e.g., Ipopt [28]), the task of solving a polynomial optimization problem (POP) to global optimality is NP-hard in general. Over the last decades, the moment-sums of squares (moment-SOS) hierarchy consisting of a sequence of increasingly tight SDP relaxations, initially established by Lasserre [18], has become a popular tool to handle polynomial optimization. The moment-SOS hierarchy features its global convergence and finite convergence under mild conditions [23]. However, the main concern on the moment-SOS hierarchy comes from its scalability as the $d$-th step of the moment-SOS hierarchy involves a semidefinite program (SDP) of size $n^{d+1}$ where $n$ is the number of decision variables of the POP. Except the first (relaxation) step of the moment-SOS hierarchy (also known as Shor’s relaxation for quadratically constrained quadratic programs (QCQP) [25]), solving higher steps of the moment-SOS hierarchy is typically limited to small-scale POPs, at least when relying on interior-point solvers. To overcome this scalability issue, one practicable way is to exploit the structure of the POP to reduce the size of SDPs arising from the moment-SOS hierarchy. Such structures include symmetry [24], correlative sparsity [19, 29], term sparsity [31, 32, 33]. The purpose of this paper is to demonstrate that the scalability of the moment-SOS hierarchy can...
be significantly improved when appropriate sparsity patterns are accessible via a thorough numerical experiment on the AC optimal power flow (AC-OPF) problem.

**Background on the AC-OPF problem.** The AC-OPF is a fundamental problem in power systems, which has been extensively studied in recent years; for a detailed introduction and recent developments, the reader is referred to the survey [5] and references therein. One can formulate the AC-OPF problem as a POP either with real variables [5, 11] or with complex variables [16]. Nonlinear programming tools can mostly produce a locally optimal solution whose global optimality is however unknown. We also mention that ill-conditioned power systems have proved the existence of computational pitfalls in which case the related nonlinear optimization algorithm can encounter unfeasibility issues, especially in the radial parts of the networks. If we are looking precisely for such unfeasibility limit behaviors, one can often encounter this issue but rather rarely when one optimizes under realistic conditions, while taking into account the maximal injections of both producers and consumers. In our study we are in the latter situation as we focus on real data coming from PGLiB instances. There have been several well-known heuristics plugged into local optimization schemes, e.g., Newton-Raphson, to make operational OPF problems more robust to solve [27]. Alternatively, one can replace the Newton-Raphson scheme by a trust-region method (more accurate but also more CPU-demanding since information about the Hessian are required) as in [1] or a Riemannian optimization approach, which may yield better algorithmic convergence guarantees, as in [14]. However such issues are less likely to arise when using semidefinite programming based on interior-point methods, since the underlying linear system to solve is different from the one considered in local solvers such as Newton-Raphson or trust-region schemes; a comprehensive study on this phenomenon could be found in [10].

Since 2006, several convex relaxation schemes (e.g., second order cone relaxations (SOCR) [15], quadratic convex relaxations (QCR) [9], tight-and-cheap conic relaxations (TCR) [6] and semidefinite (Shor’s) relaxations (SDR) [4]) have been proposed to provide lower bounds for the AC-OPF which can be then used to certify global optimality of locally optimal solutions. In particular, SDR is equivalent to the first relaxation step of the moment-SOS hierarchy. SOCR is weaker but cheaper than SDR. It was shown that TCR offers a trade-off between SOCR and SDR in terms of optimality gap and computational cost [6], and that QCR is stronger than SOCR but neither dominates nor is dominated by SDR [9]. An experimental study in [7] showed that on average TCR dominates QCR in terms of optimality gap. While these relaxations (SOCR, QCR, TCR, SDR) could be scalable to problems of large size and prove to be tight for quite a few cases [3, 9, 10], they yield significant optimality gaps for a large number of other cases. To tackle these more challenging cases, it is then mandatory to go beyond SDR, providing tighter lower bounds. Along with this line recently in [12], Gopinath et al. certified 1% global optimality for all AC-OPF instances with up to 300 buses from the AC-OPF library PGLiB using an SDP-based bound tightening approach. Relying on the complex moment-SOS hierarchy combined with a multi-order technique, Josz and Molzahn certified 0.05% global optimality for certain 2000-bus cases on a simplified AC-OPF

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1The reader may find related results on benchmarking SOCR and QCR at https://github.com/power-grid-lib/pglib-opf/blob/master/BASELINE.md.
model [16]. To certify global optimality for challenging AC-OPF instances in the general form and with thousands of buses, the CS-TSSOS hierarchy then comes into play.

The CS-TSSOS hierarchy for large-scale POPs. The CS-TSSOS hierarchy developed in [33] by the authors is a sparsity-adapted version of the moment-SOS hierarchy targeted at large-scale POPs by simultaneously exploiting correlative sparsity (CS) and term sparsity (TS). The underlying idea is the following:

1. decomposing the system into subsystems by exploiting correlative sparsity, i.e., the fact that only a few variable products occur;
2. exploiting term sparsity, i.e., the fact that the input data only contain a few terms (by comparison with the maximal possible amount), to each subsystem to further reduce the size of SDPs.

By virtue of this two-step reduction procedure, one may obtain SDP relaxations of significantly smaller size compared to the original SDP relaxations. Moreover, global convergence of the CS-TSSOS hierarchy can be still guaranteed under certain conditions [33]. Next the main concern on the CS-TSSOS hierarchy might be how it performs when being applied to real-world large-scale POPs in terms of scalability and accuracy, which will be addressed in the present paper.

Certifying global optimality for AC-OPF instances from PGLiB. We first propose a minimal initial relaxation step for the CS-TSSOS hierarchy. Being applied to the AC-OPF problem, this initial relaxation step is tighter than SDR and is less expensive than the second relaxation step of the CS-TSSOS hierarchy. Then as the main contribution of this paper, we benchmark the CS-TSSOS hierarchy through a comprehensive numerical experiment on AC-OPF instances from the AC-OPF library PGLiB v20.07 [3] with up to tens of thousands of variables and constraints. The experimental results (see Section 6) demonstrate that the CS-TSSOS hierarchy scales well with the problem size and is able to certify global optimality (in the sense of reducing optimality gap within 1%) for many of the challenging test cases. In particular, the largest instance whose global optimality is certified beyond Shor’s relaxation involves 14398 real variables and 63577 constraints. Besides, the largest instance for which the CS-TSSOS hierarchy is able to provide a smaller optimality gap than Shor’s relaxation involves 24032 real variables and 96805 constraints. To the best of our knowledge, this is the first time in the literature that one can solve higher steps of the moment-SOS hierarchy other than Shor’s relaxation for POPs of such large sizes.

2. Notation and preliminaries

A polynomial \( f \in \mathbb{R}[x] \) can be written as \( f(x) = \sum_{\alpha \in \mathcal{A}} f_\alpha x^\alpha \) with \( \mathcal{A} \subseteq \mathbb{N}^n \), \( f_\alpha \in \mathbb{R} \), and \( x^\alpha := x_1^{\alpha_1} \cdots x_n^{\alpha_n} \). The support of \( f \) is defined by \( \text{supp}(f) := \{ \alpha \in \mathcal{A} | f_\alpha \neq 0 \} \). For \( d \in \mathbb{N} \), the set \( \mathbb{N}^n_d := (x^\alpha)_{\alpha \in \mathbb{N}^n} \) is called the standard monomial basis up to degree \( d \). For convenience we abuse notation in the sequel, and denote by \( \mathbb{N}^n \) instead of \( \mathbb{N}^n_d \) the standard monomial basis and use the exponent \( \alpha \) to represent a monomial \( x^\alpha \). With \( y = (y_\alpha)_{\alpha \in \mathbb{N}^n} \subseteq \mathbb{R} \) being a sequence indexed by \( \mathbb{N}^n \), let \( L_y : \mathbb{R}[x] \to \mathbb{R} \) be the linear functional \( f = \sum_{\alpha} f_\alpha x^\alpha \mapsto L_y(f) = \sum_{\alpha} f_\alpha y_\alpha \). For
Table 1. Notation.

| Notation | Meaning |
|----------|---------|
| \(\mathbb{N}\) | \{0, 1, 2, \ldots\} |
| \(\mathbb{R}\) (resp. \(\mathbb{C}\)) | the set of real (resp. complex) numbers |
| \([m]\) | \{1, 2, \ldots, m\} |
| \([m : n]\) | \{m, m + 1, \ldots, n\} |
| \(x = (x_1, \ldots, x_n)\) | a tuple of variables |
| \(\mathbb{R}[x]\) | the ring of real \(n\)-variate polynomials |
| \(\mathbb{N}_d^n\) | \(\{(\alpha_i)_{i=1}^n \in \mathbb{N}^n \mid \sum_{i=1}^n \alpha_i \leq d\}\) |
| \(\alpha, \beta, \gamma\) | vectors in \(\mathbb{N}_d^n\) |
| \(|\cdot|\) | the cardinality of a set |
| \(S^r\) | the set of \(r \times r\) symmetric matrices |
| \(A \succeq 0\) | \(A\) is a positive semidefinite (PSD) matrix |
| \(S^r_+\) | the set of \(r \times r\) PSD matrices |
| \(A \circ B\) | the Hadamard product of matrices \(A\) and \(B\) |
| \(G(V, E)\) | a graph \(G\) with nodes \(V\) and edges \(E\) |
| \(V(G)\) (resp. \(E(G)\)) | the node (resp. edge) set of a graph \(G\) |
| \(G'\) | a chordal extension of a graph \(G\) |
| \(B_G\) | the adjacency matrix of a graph \(G\) |

\(\alpha \in \mathbb{N}^n, \mathcal{A}, \mathcal{B} \subseteq \mathbb{N}^n\), let \(\alpha + \mathcal{A} := \{\alpha + \beta \mid \beta \in \mathcal{A}\}\) and \(\mathcal{A} + \mathcal{B} := \{\alpha + \beta \mid \alpha \in \mathcal{A}, \beta \in \mathcal{B}\}\). For \(\beta = (\beta_i) \in \mathbb{N}^n\), let \(\text{supp}(\beta) := \{i \in [n] \mid \beta_i \neq 0\}\).

An (undirected) graph \(G(V, E)\), or simply \(G\), consists of a set of nodes \(V\) and a set of edges \(E \subseteq \{\{u, v\} \mid u \neq v, (u, v) \in V \times V\}\). The adjacency matrix of a graph \(G\) is denoted by \(B_G\) for which we put ones on its diagonal. A clique of a graph is a subset of nodes that induces a complete subgraph. A maximal clique is a clique that is not contained in any other clique. By definition, a chordal graph is a graph in which any cycle of length at least four has a chord\(^2\). Any non-chordal graph \(G(V, E)\) can be always extended to a chordal graph \(G'(V, E')\) by adding appropriate edges to \(E\), which is called a chordal extension of \(G(V, E)\). The chordal extension of \(G\) is usually not unique and the symbol \(G'\) is used to represent any specific chordal extension of \(G\) throughout the paper.

Given a graph \(G(V, E)\), a symmetric matrix \(Q\) with rows and columns indexed by \(V\) is said to have sparsity pattern \(G\) if \(Q_{uv} = Q_{vu} = 0\) whenever \(u \neq v\) and \(\{u, v\} \notin E\). Let \(S_G\) be the set of symmetric matrices with sparsity pattern \(G\) and let \(\Pi_G\) be the projection from \(S^{|V|}\) to the subspace \(S_G\), i.e., for \(Q \in S^{|V|}\),

\[
[\Pi_G(Q)]_{uv} = \begin{cases} Q_{uv}, & \text{if } u = v \text{ or } \{u, v\} \in E, \\ 0, & \text{otherwise}. \end{cases}
\]

The set \(\Pi_G(S^{|V|}_+)\) denotes matrices in \(S_G\) that have a PSD completion in the sense that diagonal entries, and off-diagonal entries corresponding to edges of \(G\) are

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\(^2\)A chord is an edge that joins two nonconsecutive nodes in a cycle.
fixed; other off-diagonal entries are free. More precisely, $\Pi_G(S_+^{[V]}) = \{\Pi_G(Q) \mid Q \in S_+^{[V]}\}$. For a chordal graph $G$, the following theorem due to Grone et al. gives a characterization of matrices in the PSD completable cone $\Pi_G(S_+^{[V]})$, which plays a crucial role in sparse semidefinite programming.

**Theorem 2.1** ([13], Theorem 7). Let $G(V, E)$ be a chordal graph and assume that $C_1, \ldots, C_t$ are the list of maximal cliques of $G(V, E)$. Then a matrix $Q \in \Pi_G(S_+^{[V]})$ if and only if $Q[C_i] \succeq 0$ for $i = 1, \ldots, t$, where $Q[C_i]$ denotes the principal submatrix of $Q$ indexed by the clique $C_i$.

3. THE CS-TSSOS HIERARCHY

The moment-SOS hierarchy [18] provides a sequence of increasingly tighter SDP relaxations for the following polynomial optimization problem:

\[
\begin{align*}
\inf_{x \in \mathbb{R}^n} & \quad f(x) \\
\text{s.t.} & \quad g_j(x) \geq 0, \quad j \in [m], \\
& \quad g_j(x) = 0, \quad j \in [m+1 : m+\ell],
\end{align*}
\]

where $f, g_1, \ldots, g_m \in \mathbb{R}[x]$ are all polynomials.

To state the moment hierarchy\(^3\), recall that for a given $d \in \mathbb{N}$, the $d$-th order moment matrix $M_d(y)$ associated with $y = (y_\alpha)_{\alpha \in \mathbb{N}_d}$ is defined by $[M_d(y)]_{\beta \gamma} := L_y(x^{\beta}x^{\gamma}) = y_{\beta+\gamma}, \forall \beta, \gamma \in \mathbb{N}_d$, and the $d$-th order localizing matrix $M_d(g(y))$ associated with $y$ and $g = \sum_\alpha g_\alpha x^{\alpha} \in \mathbb{R}[x]$ is defined by $[M_d(g(y))]_{\beta \gamma} := L_y(gx^{\beta}x^{\gamma}) = \sum_\alpha g_\alpha y_{\alpha+\beta+\gamma}, \forall \beta, \gamma \in \mathbb{N}_d$. Let $d_j := \lfloor \deg(g_j)/2 \rfloor$ for $j = 1, \ldots, m+\ell$ and $d_{\min} := \max \{\lfloor \deg(f)/2 \rfloor, d_1, \ldots, d_{m+\ell}\}$. Then for an integer $d \geq d_{\min}$, the $d$-th order moment relaxation for POP (3.1) is given by

\[
\begin{align*}
\inf_{y} & \quad L_y(f) \\
\text{s.t.} & \quad M_d(y) \succeq 0, \\
& \quad M_{d-d_j}(g_jy) \succeq 0, \quad j \in [m], \\
& \quad M_{d-d_j}(g_jy) = 0, \quad j \in [m+1 : m+\ell], \\
& \quad y_0 = 1.
\end{align*}
\]

We call (3.2) the dense moment hierarchy for POP (3.1), whose optima converge to the global optimum of (3.1) under mild conditions (slightly stronger than compactness of the feasible set) [18]. Unfortunately, when the relaxation order $d$ is greater than 1, the dense moment hierarchy encounters a severe scalability issue as the maximal size of PSD constraints is a combinatorial number in terms of $n$ and $d$. Therefore in the following subsections, we briefly revisit the framework of

\(^3\)We mainly focus on the moment hierarchy. The SOS hierarchy consists of the dual SDPs.
exploiting sparsity to derive a sparse moment hierarchy of remarkably smaller size for POP (3.1) in the presence of appropriate sparsity patterns. For details, the interested reader may refer to the early work on correlative sparsity by Waki et al. [19, 29] and the recent work on term sparsity by the authors [30, 31, 32, 33].

3.1. Correlative sparsity (CS). Let us from now on fix a relaxation order \( d \). By exploiting correlative sparsity, we decompose the set of variables into a tuple of subsets and then the initial system splits into a tuple of subsystems. To this end, we define the correlative sparsity pattern (csp) graph\(^4\) associated with POP (3.1) to be the graph \( G^{csp} \) with nodes \( V = [n] \) and edges \( E \) satisfying \( \{i, j\} \in E \) if one of the following conditions holds:

(i) there exists \( \alpha \in \text{supp}(f) \cup \bigcup_{k \in J^c} \text{supp}(g_k) \) such that \( \{i, j\} \subseteq \text{supp}(\alpha) \);

(ii) there exists \( k \in [m + l] \setminus (J' \cup K') \) such that \( \{i, j\} \subseteq \bigcup_{\alpha \in \text{supp}(g_k)} \text{supp}(\alpha) \),

where \( J' := \{k \in [m] \mid d_k = d\} \) and \( K' := \{k \in [m + 1 : m + l] \mid d_k = d\} \).

Let \( (G^{csp})' \) be a chordal extension of \( G^{csp} \) and \( \{I_k\}_{k \in [p]} \) be the list of maximal cliques of \( (G^{csp})' \) with \( n_k := |I_k| \). We then partition the polynomials \( g_j, j \in [m] \setminus J' \) into groups \( \{g_j \mid j \in J_k\}, k \in [p] \) which satisfy

(i) \( J_1, \ldots, J_p \subseteq [m] \setminus J' \) are pairwise disjoint and \( \bigcup_{j=1}^p J_k = [m] \setminus J' \);

(ii) for any \( j \in J_k \), \( \bigcup_{\alpha \in \text{supp}(g_j)} \text{supp}(\alpha) \subseteq I_k, k \in [p] \).

Similarly, we also partition the polynomials \( g_j, j \in [m + 1 : m + l] \setminus K' \) into groups
\( \{g_j \mid j \in K_k\}, k \in [p] \).

For any \( k \in [p] \), let \( M_d(y, I_k) \) (resp. \( M_d(gy, I_k) \)) be the moment (resp. localizing) submatrix obtained from \( M_d(y) \) (resp. \( M_d(gy) \)) by retaining only those rows and columns indexed by \( \beta \in \mathbb{N}_n^d \) of \( M_d(y) \) (resp. \( M_d(gy) \)) with \( \text{supp}(\beta) \subseteq I_k \).

Example 3.1. Consider the POP:

\[
\begin{aligned}
\inf_{x \in \mathbb{R}^3} & \quad x_1^2 x_2 + x_2 x_3^2 \\
\text{s.t.} & \quad 1 - x_1^2 - x_2^2 \geq 0, \\
& \quad 1 - x_2^2 - x_3^2 \geq 0, \\
& \quad x_1^4 + x_2 x_3 = 1.
\end{aligned}
\]

Let us take the relaxation order \( d = 2 \). Then the csp graph is shown in Figure 2, which contains two maximal cliques: \( \{x_1, x_2\} \) and \( \{x_2, x_3\} \).

\[\text{Figure 2. Illustration for correlative sparsity.}\]

\(^4\)We adopt the idea of “monomial sparsity” introduced in [16] for the definition of csp graphs, which thus is slightly different from the original definition given in [29].
3.2. Term sparsity (TS). We next apply an iterative procedure to exploit term sparsity for each subsystem involving variables $x(I_k) := \{ x_i \mid i \in I_k \}$ for $k \in [p]$. The intuition behind this procedure is the following: starting with a minimal initial set of moments, we expand the set of moments that is taken into account in the moment relaxation by iteratively performing chordal extension to the related graphs inspired by Theorem 2.1. More concretely, let $\mathcal{A} := \text{supp}(f) \cup \bigcup_{j=1}^{m} \text{supp}(g_j)$ and $\mathcal{A}_k := \{ \alpha \in \mathcal{A} \mid \text{supp}(\alpha) \subseteq I_k \}$ for $k \in [p]$. We define $G_{d,k,0}^{(0)}$ to be the graph with nodes $V_{d,k,0} := \mathbb{N}^{n_k}$ and edges

$$E(G_{d,k,0}^{(0)}) := \{ \{ \beta, \gamma \} \mid \beta, \gamma \in V_{d,k,0}, \beta + \gamma \in \mathcal{A}_k \cup (2\mathbb{N})^n \}.$$  

Note that here we embed $\mathbb{N}^{n_k}$ into $\mathbb{N}$ by specifying the $i$-th coordinate to be zero when $i \in [n] \setminus I_k$.

Example 3.2. Consider again the POP in Example 3.1 with the relaxation order $d = 2$. Since there are two variable cliques derived from the correlative sparsity pattern, we have $p = 2$. Figure 3 illustrates the term sparsity pattern of this POP.

![Figure 3](image)

For the sake of convenience, we set $g_0 := 1$ and $d_0 := 0$ hereafter and for a graph $G(V, E)$ with $V \subseteq \mathbb{N}$, let $\text{supp}(G) := \{ \beta + \gamma \mid \beta = \gamma \in V \text{ or } \{ \beta, \gamma \} \in E \}$. Assume that $G_{d,k,0}^{(0)}, j \in J_k \cup K_k, k \in [p]$ are all empty graphs with $V_{d,k,0} := \mathbb{N}^{n_k}$. Now for each $j \in \{ 0 \} \cup J_k \cup K_k, k \in [p]$, we iteratively define an ascending chain of graphs $(G_{d,k,0}^{(s)})_{s \geq 1}$ by

$$G_{d,k,0}^{(s)} := (F_{d,k,0}^{(s)})',$$

where $F_{d,k,0}^{(s)}$ is the graph with nodes $V_{d,k,0}$ and edges

$$E(F_{d,k,0}^{(s)}) = \{ \{ \beta, \gamma \} \mid \beta, \gamma \in V_{d,k,0}, (\beta + \gamma + \text{supp}(g_j)) \cap \mathcal{C}_d^{(s-1)} \neq \emptyset \},$$

with

$$\mathcal{C}_d^{(s-1)} := \bigcup_{k=1}^{p} \bigcup_{j \in \{ 0 \} \cup J_k \cup K_k} \left( \text{supp}(g_j) + \text{supp}(G_{d,k,0}^{(s-1)}) \right).$$
Let $r_{d,k,j} := |\mathbb{N}^n_{d,k,j}| = (m_k + 1 - d_j) - d_j$ for all $k, j$. Then for each $s \geq 1$, the moment relaxation based on correlative-term sparsity for POP (3.1) is given by

\[
\begin{align*}
\inf_{y} & \quad L_y(f) \\
\text{s.t.} & \quad B_{G_{d,k,0}^{(s)}} \circ M_{d}(y, I_k) \in \Pi_{G_{d,k,0}^{(s)}} (S_r^{d,k,0}), \quad k \in [p], \\
& \quad B_{G_{d,k,j}^{(s)}} \circ M_{d-d_j}(g_j y, I_k) \in \Pi_{G_{d,k,j}^{(s)}} (S_r^{d,k,j}), \quad j \in J_k, k \in [p], \\
& \quad B_{G_{d,k,i}^{(s)}} \circ M_{d-d_j}(g_j y, I_k) = 0, \quad j \in K_k, k \in [p], \\
& \quad L_y(g_j) \geq 0, \quad j \in J', \\
& \quad L_y(g_j) = 0, \quad j \in K', \\
& \quad y_0 = 1.
\end{align*}
\]

The above hierarchy is called the CS-TSSOS hierarchy, which is indexed by two parameters: the relaxation order $d$ and the sparse order $s$.

3.3. The minimal initial relaxation step. For POP (3.1), suppose that $f$ is not a homogeneous polynomial or the polynomials $g_j, j \in [m+1]$ are of different degrees as in the case of the AC-OPF problem. Then instead of using the uniform minimum relaxation order $d_{\text{min}}$, it might be more beneficial, from the computational point of view, to assign different relaxation orders to different subsystems obtained from the correlative sparsity pattern for the initial relaxation step of the CS-TSSOS hierarchy. To this end, we redefine the CSP graph $G^{\text{csp}}(V, E)$ as follows: $V = [n]$ and $\{i, j\} \in E$ whenever there exists $\alpha \in \text{supp}(f) \cup \bigcup_{j \in [m+1]} \text{supp}(g_j)$ such that $\{i, j\} \subseteq \text{supp}(\alpha)$. This is clearly a subgraph of $G^{\text{csp}}$ defined in Section 3.1 and hence typically admits a smaller chordal extension. Let $(G^{\text{csp}})'$ be a chordal extension of $G^{\text{csp}}$ and $\{I_k\}_{k \in [p]}$ be the list of maximal cliques of $(G^{\text{csp}})'$ with $n_k := |I_k|$. Now we partition the polynomials $g_j, j \in [m]$ into groups $\{g_j | j \in J_k\}_{k \in [p]}$ and $\{g_j | j \in J'\}$ which satisfy

(i) $J_1, \ldots, J_p, J' \subseteq [m]$ are pairwise disjoint and $\bigcup_{k=1}^p J_k \cup J' = [m]$;
(ii) for any $j \in J_k, \bigcup_{\alpha \in \text{supp}(g_j)} \text{supp}(\alpha) \subseteq I_k, k \in [p]$;
(iii) for any $j \in J', \bigcup_{\alpha \in \text{supp}(g_j)} \text{supp}(\alpha) \not\subseteq I_k$ for all $k \in [p]$.

Similarly, we also partition the polynomials $g_j, j \in [m + 1 : m + l]$ into groups $\{g_j | j \in K_k\}_{k \in [p]}$ and $\{g_j | j \in K'\}$.

Assume that $f$ decomposes as $f = \sum_{k \in [p]} f_k$ such that $\bigcup_{\alpha \in \text{supp}(f_k)} \text{supp}(\alpha) \subseteq I_k$ for $k \in [p]$. We define the vector of minimum relaxation orders $\mathbf{o} = (o_k)_{k \in [p]}$ with $o_k := \max \{\{d_j : j \in J_k \cup K_k\} \cup \{\lceil \deg(f_k)/2 \rceil\}\}$. Then with $s \geq 1$, we define
the following minimal initial relaxation step for the CS-TSSOS hierarchy:

\[
\begin{align*}
\inf_{y} \quad & L_y(f) \\
\text{s.t.} \quad & B_{G_{ok,k,0}}^{(s)} \circ M_{ok}(y, I_k) \in \Pi_{G_{ok,k,0}}(S_{+}^{t_{k,0}}), \quad k \in [p], \\
& M_1(y, I_k) \geq 0, \quad k \in [p], \\
& B_{G_{ok,k,j}}^{(s)} \circ M_{ok-d_j}(g_j y, I_k) \in \Pi_{G_{ok,k,j}}(S_{+}^{t_{k,j}}), \quad j \in J_k, k \in [p], \\
& L_y(g_j) \geq 0, \quad j \in J', \\
& B_{G_{ok,k,j}}^{(s)} \circ M_{ok-d_j}(g_j y, I_k) = 0, \quad j \in K_k, k \in [p], \\
& L_y(g_j) = 0, \quad j \in R', \\
& y_0 = 1,
\end{align*}
\]  

(3.8)

where \( G_{ok,k,j}^{(s)} \), \( j \in J_k \cup K_k, k \in [p] \) are defined as in Section 3.2 and \( t_{k,j} := \binom{n_k + o_k - d_j}{o_k - d_j} \) for all \( k, j \). Note that in (3.8) we add the PSD constraint on each first-order moment matrix \( M_1(y, I_k) \) to strengthen the relaxation.

The CS-TSSOS hierarchy is implemented in the Julia package TSSOS\(^5\). In TSSOS, the minimal initial relaxation step is accessible via the commands \texttt{cs\_tssos\_first} and \texttt{cs\_tssos\_higher}! by setting the relaxation order to be "\texttt{\textasciitilde}min". For an introduction to TSSOS, the reader is referred to [22].

4. Problem formulation of AC-OPF

The AC-OPF problem aims to minimize the generation cost of an alternating current transmission network under the physical constraints (Kirchhoff’s laws, Ohm’s law) as well as operational constraints, which can be formulated as the following POP in complex variables:

\[
\begin{align*}
\inf_{V_i, S_i^R, S_i^I} \quad & \sum_{k \in G} (c_{2k}(\Re(S_k^R))^2 + c_{1k}\Re(S_k^R) + c_{ok}) \\
\text{s.t.} \quad & \angle V_i = 0, \\
& S_k^R \leq S_k^q \leq S_k^{qu}, \quad \forall k \in G, \\
& \nu_i \leq |V_i| \leq \nu_i, \quad \forall i \in N, \\
& \sum_{k \in G} S_k^R - S_i^R - Y_i^\ast |V_i|^2 = \sum_{(i,j) \in E_i} S_{ij}, \quad \forall i \in N, \\
& S_{ij} = (Y_{ij} - \frac{1}{2}\nu_i^2 |V_i|^2 - Y_{ij}^\ast V_i^V_{ij}), \quad \forall (i, j) \in E, \\
& S_{ij} = (Y_{ij} - \frac{1}{2}\nu_i^2 |V_i|^2 - Y_{ij}^\ast V_i^V_{ij}), \quad \forall (i, j) \in E, \\
& |S_{ij}| \leq S_i^q, \quad \forall (i, j) \in E \cup R, \\
& \theta_{ij}^{\Delta L} \leq \angle(V_i V_j) \leq \theta_{ij}^{\Delta U}, \quad \forall (i, j) \in E.
\end{align*}
\]  

(4.1)

The meaning of the symbols in (4.1) is as follows: \( N \) - the set of buses, \( G \) - the set of generators, \( G_i \) - the set of generators connected to bus \( i \), \( E \) - the set of \texttt{from} branches, \( E_i \) and \( E_i^R \) - the subsets of branches that are incident to bus \( i \), \( i \) - imaginary unit, \( V_i \) - the voltage at bus \( i \), \( S_k^R \) - the power generation at generator \( k \), \( S_{ij} \) - the power flow from bus \( i \) to bus \( j \), \( \Re(\cdot) \) - real part of a complex number, \( \angle(\cdot) \) - angle of a complex number, \( |\cdot| \) - magnitude of a complex number, \( (\cdot)^\ast \) - conjugate of a complex number, \( r \) - the voltage angle reference.

\(^5\)TSSOS is freely available at https://github.com/wangjie212/TSSOS.
bus. All symbols in boldface are constants \((c_0 k, c_1 k, c_2 k, u^l_i, u^u_i, s^l_{ij}, s^u_{ij}, \theta^l_{ij}, \theta^u_{ij} \in \mathbb{R}, S^l_{ij}, S^u_{ij}, S^d_i, Y_s, b_c, T_{ij} \in \mathbb{C})\). For a full description on the AC-OPF problem, the reader may refer to \([3]\). By introducing real variables for both real and imaginary parts of each complex variable, we can convert the AC-OPF problem to a POP involving only real variables\(^6\).

To tackle an AC-OPF instance, we first compute a locally optimal solution with a local solver (e.g., \texttt{Ipopt} \([28]\)) and then rely on lower bounds obtained from certain relaxation schemes (SOCR/QCR/TCR/SDR/CS-TSSOS) to certify 1\% global optimality. Suppose that the optimum reported by the local solver is \(AC\) and the lower bound given by a certain convex relaxation is \(\text{opt}\). Then the \textit{optimality gap} is defined by

\[
\text{gap} := \frac{AC - \text{opt}}{AC} \times 100\%.
\]

As in \([12]\), if the optimality gap is less than 1\%, then we accept the locally optimal solution to be globally optimal. The procedure here for certifying global optimality of solutions of an AC-OPF instance is summarized in Algorithm 1.

\begin{algorithm}
\caption{Algorithm 1}
\begin{algorithmic}
\STATE \textbf{Input:} Data for an AC-OPF instance of (4.1)
\STATE \textbf{Output:} A local solution and the optimality gap\(^1\)
\STATE 1: Build the POP model (3.1) with the given data;
\STATE 2: Compute a local solution (\(\text{sol}\)) of the POP model with a local solver, and denote the optimum by \(AC\);
\STATE 3: Build the SDP relaxation (3.8) for the POP model;
\STATE 4: Solve the SDP with \texttt{Mosek}, and denote the optimum by \(\text{opt}\);
\STATE 5: \(\text{gap} \leftarrow \frac{AC - \text{opt}}{AC} \times 100\%\);
\STATE 6: \textbf{return} \text{sol, gap}
\end{algorithmic}
\end{algorithm}

In our experiments, we eliminate the power flow variables \(S_{ij}\) from (4.1) so that it only involves the voltage variable \(V_i\) and the power generation variables \(S^l_{k}\). This reformulation is crucial to derive a CS-TSSOS hierarchy of lower complexity, which yet introduces a quartic constraint corresponding to \(S_{ij} S^*_{ij} \leq (s^l_{ij})^2\). To implement the first order relaxation, we then have to relax this quartic constraint to a quadratic constraint using the trick described in \([5, \text{Sec. 5.3}]\). The minimal initial relaxation step of the CS-TSSOS hierarchy for (4.1) is able to provide a tighter lower bound than the first order relaxation and is less expensive than the second order relaxation. Therefore, we hereafter refer to it as the 1.5th order relaxation.

5. Experimental settings

\textbf{Challenging test cases.} Our test cases are selected from the AC-OPF library PGLiB v20.07 which provides various AC-OPF instances for benchmarking AC-OPF algorithms. For an introduction to this library, the reader is referred to \([3]\). We observe that for a number of instances in PGLiB, the SOCR approach is able to close the gap to below 1\% and these instances are not particularly interesting as our purpose is to certify 1\% global optimality for more challenging cases. To that

\footnote{The expressions involving angles of complex variables can be converted to polynomials by using \(\tan(\angle z) = y/x\) for \(z = x + iy \in \mathbb{C}\).}
end, we select test cases from PGLiB (with no more than 25000 buses) for which SOCR yields an optimality gap greater than 1%. There are 115 such instances in total. For each instance, with TSSOS we initially solve the first order relaxation and if this relaxation fails to certify 1% global optimality, we further solve the 1.5th order relaxation with $s = 1$. Here Mosek 9.0 [2] is employed as an SDP solver with the default settings.

**Chordal extension.** To achieve a good balance between the computational cost and the approximation quality of lower bounds, two types of chordal extensions are used in the computation. For correlative sparsity, we use approximately smallest chordal extensions which give rise to small clique numbers. For term sparsity, instead we use maximal chordal extensions which make every connected component to be a complete subgraph by setting $TS = "block"$ in TSSOS.

![Figure 4](image.png)

**Figure 4.** Illustration for smallest chordal extension (left) and maximal chordal extension (right): the dashed edges are added via chordal extension.

**Scaling of polynomial coefficients.** To improve the numerical conditioning of the SDP relaxations, we scale the coefficients of $f$ and $g_j$ so that they lie in the interval $[-1, 1]$ before building the SDP relaxations.

**Computational resources.** Instances with no more than 3500 buses (except 2853_sdet and 2869_pegase) were computed on a laptop with an Intel Core i5-8265U@1.60GHz CPU and 8GB RAM memory; instances with more than 3500 buses (including 2853_sdet and 2869_pegase) were computed on a server with an Intel Xeon E5-2695v4@2.10GHz CPU and 128GB RAM memory.

**Table 2.** Notations for the numerical results.

| Notation | Meaning |
|----------|---------|
| AC       | local optimum (available from PGLiB) |
| mc       | maximal size of variable cliques |
| mb       | maximal size of SDP blocks |
| opt      | optimum of SDP relaxations |
| time     | running time in seconds |
| gap      | optimality gap (%) |
| *        | encountering a numerical error |
| -        | out of memory |
6. Computational results and discussion

The computational results are summarized in Table 3–5 corresponding to three operational conditions, denoted by “typical”, “congested” and “small angle differences”, respectively, where the timing includes the time for pre-processing (to obtain the block structure), the time for building SDP and the time for solving SDP. Note that the maximal size of variable cliques varies from 6 to 218 among the tested cases. According to the tables, we can draw the following conclusions.

Reducing the optimality gap. As we would expect, the 1.5th order relaxation provides tighter lower bounds than the first order relaxation. Indeed, when it is solvable, the 1.5th order relaxation always reduces the optimality gap (unless the lower bounds given by the first order relaxation are already globally optimal). The largest instance for which the 1.5th order relaxation is solvable is 10000goc and its corresponding POP involves 24032 real variables and 96805 constraints. The improvement of optimality gaps with the 1.5th order relaxation is significant on quite a few cases. For instance, the first order relaxation yields a optimality gap of 42.96\% on 30 as under congested operating conditions while the 1.5th order relaxation yields a optimality gap of merely 0.01\%. On the other hand, as the cost of these improvements solving the 1.5th order relaxation typically spends significantly more time than solving the first order relaxation.

Certifying 1% global optimality. The first order relaxation is able to certify 1% global optimality for 29 out of all 115 instances. The 1.5th order relaxation is able to certify 1% global optimality for 29 out of the remaining 86 instances. The largest instance for which we are able to certify 1% global optimality with the 1.5th order relaxation is 6515rte and its corresponding POP involves 14398 real variables and 63577 constraints. One would expect that solving the second order relaxation could certify global optimality for more instances. However this is too expensive to implement for large-scale instances in practice.

Computational burden. The computational burden of the CS-TSSOS relaxations heavily relies on the maximal size of variable cliques. This is because large
variable cliques usually lead to SDP matrices of large size in the resulting CS-TSSOS relaxations. It takes 63785 seconds to solve the 1.5th order relaxation for the case 4020_goc under congested operating conditions as it involves a variable clique of size 120. For similar reasons, Mosek runs out of memory with the 1.5th order relaxation for the cases 9241_pegase, 9591_goc, 10480_goc, 13659_pegase, 19402_goc, 24464_goc.

**Numerical issues.** Even though we have scaled polynomial coefficients to improve numerical conditioning of the resulting SDPs, we observe that in numerous cases (especially when solving the 1.5th order relaxation), the termination status of Mosek is “slow progress”, which means that Mosek does not converge to the default tolerance although the solver usually still gives a fairly good near-optimal solution in this case. Moreover, there are 12 even more challenging instances for which Mosek encounters severe numerical issues with the 1.5th order relaxation and fails in converging to the optimum. This indicates that there is still room for improvement in order to tackle these challenging SDPs.

### Table 3. Results for AC-OPF problems: typical operating conditions.

| case name | AC | first order | 1.5th order |
|-----------|----|-------------|-------------|
|           | opt | time | mb | gap | opt | time | mb | gap |
| 3_lmbd    | 5.8126e3 | 5.7455e3 | 0.10 | 5 | 1.15 | 6 | 5.8126e3 | 0.12 | 22 | **0.00** |
| 5_pjm     | 1.7552e4 | 1.4997e4 | 0.15 | 6 | 14.56 | 6 | 1.7552e4 | 0.58 | 22 | **0.10** |
| 30_joc    | 8.2085e3 | 7.5472e3 | 0.22 | 8 | 8.06 | 8 | 8.2085e3 | 0.99 | 22 | **0.00** |
| 162_jec_ge | 1.0164e5 | 2.15 | 28 | 5.96 | 28 | 1.0164e5 | 99.1 | 74 | 1.51 |
| 240_psrc  | 3.3297e6 | 3.2512e6 | 2.39 | 16 | 2.36 | 16 | 3.3297e6 | 28.6 | 44 | **0.64** |
| 300_jec   | 5.6522e3 | 5.5423e3 | 0.12 | 5 | 0.00 | 5 | 5.6522e3 | 0.58 | 22 | **0.00** |
| 588_sdet  | 3.3134e5 | 3.0886e5 | 4.37 | 18 | 1.37 | 18 | 3.3134e5 | 99.1 | 74 | 1.51 |
| 793_goc   | 2.6020e5 | 2.5636e5 | 5.35 | 18 | 1.47 | 18 | 2.6020e5 | 66.1 | 33 | **0.34** |
| 1888_dte  | 1.4025e6 | 1.3666e6 | 30.0 | 26 | 2.56 | 26 | 1.4025e6 | 458 | 56 | 1.92 |
| 2312_goc  | 4.133e6  | 3.4335e6 | 87.8 | 68 | 1.58 | 68 | 4.133e6  | 997 | 81 | **0.62** |
| 2338800   | 1.8682e6 | 1.8682e6 | 63.0 | 50 | 0.52 | 48 | 1.8682e6 | 945 | 77 | **0.19** |
| 2742_goc  | 2.7571e6 | 2.7561e6 | 703 | 92 | 0.04 | 92 | 2.7571e6 | 1886 | 76 | 0.00 |
| 2886_pgas  | 2.4624e6 | 2.4384e6 | 85.0 | 26 | 0.97 | 26 | 2.4624e6 | 3641 | 191 | **0.22** |
| 3016_lmbd  | 2.6008e6 | 2.5828e6 | 123 | 52 | 0.69 | 52 | 2.6008e6 | 1969 | 81 | **0.23** |
| 3022_goc  | 6.0138e5 | 5.9277e5 | 115 | 48 | 1.43 | 50 | 5.9277e5 | 1886 | 76 | **0.47** |
| 4020_goc  | 8.2225e6 | 8.2208e6 | 2556 | 112 | 0.02 | 112 | 8.2225e6 | 997 | 81 | **0.61** |
| 4661_sdet  | 2.2513e6 | 2.2246e6 | 25746 | 204 | 1.18 | 218 | * | * | 285 | * |
| 4917_goc  | 1.3878e6 | 1.3658e6 | 276 | 64 | 1.59 | 68 | 1.3878e6 | 29562 | 110 | **0.61** |
| 6408_dte  | 2.0697e6 | 2.0546e6 | 415 | 54 | 0.73 | 54 | 2.0546e6 | 285 | 81 | **0.41** |
| 6470_dte  | 2.2376e6 | 2.2060e6 | 478 | 54 | 1.41 | 58 | * | * | 98 | * |
| 6495_dte  | 3.0678e6 | 2.6327e6 | 426 | 56 | 14.18 | 54 | * | * | 108 | * |
| 6515_dte  | 2.8255e6 | 2.6536e6 | 460 | 56 | 5.99 | 54 | * | * | 108 | * |
| 9241_pgas  | 6.2431e6 | 6.1330e6 | 982 | 64 | 1.76 | 64 | - | - | 1268 | - |
| 10000_goc | 1.3540e6 | 1.3460e6 | 1714 | 84 | 0.59 | 84 | 1.3540e6 | 228 | 81 | **0.19** |
| 10480_goc | 2.3146e6 | 2.3051e6 | 8559 | 136 | 0.41 | 136 | 2.3146e6 | 8559 | 136 | **0.41** |
| 13659_pgas | 8.9480e6 | 8.8707e6 | 1808 | 64 | 0.86 | 64 | 8.9480e6 | 1808 | 64 | **0.86** |
| 19402_goc | 1.9778e6 | 1.9752e6 | 37157 | 180 | 0.13 | 180 | 1.9752e6 | 37157 | 180 | **0.13** |
Table 4. Results for AC-OPF problems: congested operating conditions.

| case name | AC | first order | 1.5th order |
|-----------|----|-------------|-------------|
|           | opt | time | mb | gap | mc | opt | time | mb | gap |
| 3_inbd    | 1.1236e4 | 34.995 | 6 | 1.1236e4 | 0.23 | 22 | 0.00 |
| 5_pjm     | 7.6377e4 | 34.995 | 6 | 7.6377e4 | 0.55 | 22 | 0.00 |
| 14_lee    | 5.9994e3 | 5.6886e3 | 6 | 5.9994e3 | 0.54 | 22 | 0.00 |
| 21_lee_crs | 1.3495e5 | 1.2639e5 | 10 | 1.3495e5 | 1.52 | 31 | 0.75 |
| 30_sp     | 4.9996e3 | 2.8492e3 | 8 | 4.9996e3 | 2.41 | 22 | 0.01 |
| 30_lee    | 1.8041e4 | 1.7253e4 | 8 | 1.8041e4 | 1.24 | 22 | 0.00 |
| 39_epi    | 2.0967e5 | 2.4226e5 | 12 | 2.4226e5 | 2.72 | 25 | 0.00 |
| 73_lee_crs | 4.2263e5 | 3.9912e5 | 12 | 4.2263e5 | 6.77 | 30 | 1.82 |
| 89_pegas  | 1.2718e5 | 1.0053e5 | 24 | 1.0053e5 | 1.0188e5 | 1404 | 20.29 |
| 118_lee   | 2.4224e5 | 1.9375e5 | 24 | 2.4224e5 | 11.5 | 37 | 8.56 |
| 162_lee_lte | 1.2099e5 | 1.1206e5 | 28 | 1.2099e5 | 84.1 | 74 | 1.19 |
| 170_goc   | 1.9320e6 | 1.7224e6 | 10 | 1.9226e6 | 9.69 | 37 | 0.48 |
| 500_goc   | 6.9241e5 | 6.6004e5 | 18 | 6.7825e5 | 78.0 | 50 | 2.05 |
| 588_sdet  | 3.9476e5 | 3.9026e5 | 18 | 3.9414e5 | 57.0 | 32 | 0.15 |
| 793_goc   | 3.1885e5 | 2.9796e5 | 18 | 3.1885e5 | 79.2 | 33 | 1.56 |
| 2000_goc  | 1.4686e6 | 1.4147e6 | 42 | 1.4686e6 | 1094 | 62 | 0.51 |
| 2312_goc  | 5.7152e5 | 4.7872e5 | 93.4 | 5.7152e5 | 972 | 81 | 7.77 |
| 2738_goc  | 6.5394e5 | 5.8042e5 | 89.6 | 5.8042e5 | 2619 | 90 | 2.25 |
| 2738_goc  | 3.6513e5 | 3.4557e5 | 48 | 3.4557e5 | 1653 | 77 | 5.40 |
| 2794_goc  | 6.4219e5 | 5.0824e5 | 772 | 6.4219e5 | 4644 | 108 | 5.45 |
| 2854_sdet | 2.4578e6 | 2.3869e6 | 118 | 2.4578e6 | 1029 | 293 | 0.54 |
| 2869_pegas | 2.9858e6 | 2.9604e6 | 90.2 | 2.9858e6 | 5409 | 191 | 0.35 |
| 3022_goc  | 6.5185e5 | 6.2346e5 | 102 | 6.4076e5 | 1519 | 76 | 1.71 |
| 3120_goc  | 9.3599e5 | 7.6012e5 | 138 | 7.6012e5 | 1627 | 70 | 8.93 |
| 337 goc   | 5.8406e6 | 5.5376e6 | 222 | 5.5376e6 | 2619 | 90 | 2.25 |
| 3970_goc  | 1.4241e6 | 1.0087e6 | 2409 | 1.0087e6 | 15482 | 135 | 24.73 |
| 4029_goc  | 1.2970e6 | 1.0836e6 | 3523 | 1.0836e6 | 63785 | 174 | 13.57 |
| 4601_goc  | 7.9235e6 | 6.7525e6 | 2143 | 6.7525e6 | 17249 | 125 | 6.74 |
| 4619_goc  | 1.0299e6 | 9.6351e5 | 1782 | 9.6351e5 | 18348 | 132 | 3.13 |
| 4661_goc  | 2.0953e6 | 2.6112e6 | 15822 | 2.6112e6 | 218 | 285 | 2.85 |
| 4837_goc  | 1.1578e6 | 1.0769e6 | 500 | 1.0769e6 | 8723 | 132 | 5.45 |
| 4917_goc  | 1.5479e6 | 1.4670e6 | 259 | 1.4670e6 | 4688 | 110 | 1.93 |
| 6470_goc  | 2.6005e6 | 2.5795e6 | 427 | 2.5795e6 | 98 | 98 | 0.59 |
| 6495_goc  | 2.9750e6 | 2.9092e6 | 453 | 2.9092e6 | 108 | 108 | 0.86 |
| 6515_goc  | 3.0617e6 | 2.9996e6 | 421 | 2.9996e6 | 8456 | 108 | 0.60 |
| 9241_pegas | 7.0121e6 | 6.8784e6 | 865 | 6.8784e6 | - | - | 1268 |
| 9591_goc  | 1.4259e6 | 1.2425e6 | 7674 | 148 | 1286 | 134 | - | - |
| 10000_goc | 2.3728e6 | 2.1977e6 | 2564 | 2.1977e6 | 27179 | 97 | 2.20 |
| 10481_goc | 2.7627e6 | 2.6680e6 | 8791 | 2.6680e6 | - | - | 208 |
| 13659_pegas | 9.2842e6 | 9.1360e6 | 1599 | 9.1360e6 | - | - | 1268 |
| 1940_goc  | 2.3087e6 | 2.3290e6 | 32465 | 2.3290e6 | - | - | 242 |
| 24464_goc | 2.4723e6 | 2.4177e6 | 11760 | 2.4177e6 | - | - | 172 |
Table 5. Results for AC-OPF problems: small angle difference conditions.

| case name | AC    | 1st order | 5th order |
|-----------|-------|-----------|-----------|
|           | opt   | time | mb |gap| opt | time | mb |gap|
| 3_mbd     | 5.0593e3 | 5.7463e3 | 0.11 | 5 | 5.9593e3 | 0.13 | 22 | 0.00 |
| 5_pjm     | 2.6109e4 | 2.6109e4 | 0.12 | 6 | 0.00 |
| 14_ltee   | 2.7768e3 | 2.7743e3 | 0.14 | 6 | 0.09 |
| 24_ltee_2ts | 7.6918e4 | 7.3555e4 | 0.21 | 10 | 4.37 | 7.4852e4 | 1.64 | 31 | 2.69 |
| 30_ltee   | 8.9735e2 | 8.9272e2 | 0.19 | 8 | 0.23 |
| 30_ltee   | 8.2085e3 | 7.5472e3 | 0.30 | 8 | 0.00 |
| 73_ltee_2ts | 2.2760e5 | 2.2316e5 | 0.55 | 12 | 2.74 | 2.2447e5 | 5.01 | 36 | 1.38 |
| 118_ltee  | 1.0516e5 | 1.0191e5 | 0.79 | 10 | 3.10 | 1.0313e5 | 10.8 | 37 | 1.93 |
| 162_ltee_llc | 1.0890e5 | 1.0282e5 | 2.46 | 28 | 5.40 | 1.0740e5 | 105 | 74 | 1.19 |
| 179_goc   | 7.6186e5 | 7.3261e5 | 1.39 | 10 | 1.21 | 7.5573e5 | 11.8 | 37 | 0.80 |
| 240_goc   | 3.4054e6 | 3.2712e6 | 3.00 | 16 | 3.76 | 3.3128e6 | 33.8 | 44 | 2.72 |
| 300_ltee  | 5.6570e5 | 5.1662e5 | 2.73 | 16 | 0.72 | 5.6570e5 | 25.2 | 40 | 0.00 |
| 500_goc   | 4.8740e5 | 4.6043e5 | 6.52 | 18 | 5.53 | 4.6099e5 | 67.7 | 50 | 5.42 |
| 588_gdet  | 3.2936e5 | 3.1233e5 | 5.12 | 18 | 5.17 | 3.1898e5 | 56.6 | 32 | 3.15 |
| 790_goc   | 2.8540e5 | 2.7135e5 | 5.61 | 18 | 5.06 | 2.7727e5 | 76.0 | 33 | 2.98 |
| 1354 pegase_2 | 1.2588e6 | 1.2172e6 | 19.8 | 26 | 3.31 | 1.2582e6 | 387 | 49 | 0.05 |
| 1888_rte  | 1.4139e6 | 1.3666e6 | 31.2 | 26 | 3.34 | 1.3756e6 | 497 | 56 | 2.71 |
| 2000_goc  | 9.9298e5 | 9.8406e5 | 50.9 | 42 | 0.89 | 9.8435e5 | 1052 | 62 | 0.86 |
| 2312_goc  | 4.6235e5 | 4.4719e5 | 121 | 68 | 3.28 | 4.5676e5 | 1009 | 81 | 1.21 |
| 2383_wp_1k | 1.9112e6 | 1.9041e6 | 65.6 | 50 | 0.37 | 1.9060e6 | 937 | 77 | 0.27 |
| 2736_sp_1k | 1.3260e6 | 1.3229e6 | 89.5 | 50 | 0.28 |
| 2737_op_1k | 7.9095e5 | 7.8672e5 | 76.3 | 48 | 0.53 |
| 2742_goc  | 2.7571e5 | 2.7561e5 | 686 | 92 | 0.04 |
| 2746_wp_1k | 1.2337e6 | 1.2248e6 | 79.1 | 48 | 0.72 |
| 2746_wp_1k | 1.6669e6 | 1.6601e6 | 83.1 | 50 | 0.41 |
| 2854_gdet  | 2.0692e6 | 2.0303e6 | 106 | 40 | 1.88 | 2.0537e6 | 40671 | 293 | 0.75 |
| 2869_pugase | 2.4687e6 | 2.4477e6 | 85.4 | 26 | 0.85 |
| 3012_wp_1k | 2.6192e6 | 2.5994e6 | 97.1 | 52 | 0.76 |
| 3022_goc  | 6.0143e5 | 5.9278e5 | 93.4 | 48 | 1.44 | 5.9859e5 | 1340 | 76 | 0.47 |
| 3128_sp_1k | 2.1740e6 | 2.1311e6 | 117 | 52 | 0.64 |
| 4020_goc  | 8.8969e5 | 8.4328e5 | 2746 | 112 | 5.32 | 8.7038e5 | 43180 | 174 | 2.17 |
| 4601_goc  | 8.7803e5 | 8.3370e5 | 1763 | 108 | 5.05 | 8.3417e5 | 15585 | 125 | 4.96 |
| 4619_goc  | 4.8435e5 | 4.8106e5 | 1387 | 82 | 0.68 |
| 461_gdet  | 2.2610e6 | 2.2337e6 | 16144 | 204 | 1.21 | 218 | * | * | 285 | * |
| 4917_goc  | 1.3890e6 | 1.3860e6 | 260 | 64 | 1.62 | 1.3800e6 | 4914 | 110 | 0.65 |
| 6408_rte  | 2.0697e6 | 2.0546e6 | 399 | 54 | 0.73 |
| 6408_rte  | 2.2416e6 | 2.2106e6 | 451 | 54 | 1.41 | 58 | * | * | 98 | * |
| 6409_rte  | 3.0678e6 | 2.6324e6 | 404 | 56 | 14.19 | 54 | * | * | 108 | * |
| 6515_rte  | 2.8698e6 | 2.6565e6 | 399 | 56 | 7.43 | 54 | * | * | 108 | * |
| 9241_pugase | 6.3158e6 | 6.1696e6 | 912 | 64 | 2.36 | 64 | - | - | 1268 | - |
| 9591_goc  | 1.1674e6 | 1.0712e6 | 7835 | 148 | 8.24 | 134 | - | - | 201 | - |
| 10000_goc | 1.4902e6 | 1.4204e6 | 2508 | 84 | 4.68 | 1.4212e6 | 25340 | 97 | 4.63 |
| 10480_goc | 2.3147e6 | 2.3051e6 | 6522 | 136 | 0.42 |
| 13659_pugase | 9.0422e6 | 8.9142e6 | 1053 | 64 | 1.42 | 64 | - | - | 1268 | - |
| 19402_goc | 1.9838e6 | 1.9783e6 | 30122 | 180 | 0.28 |
| 24464_goc | 2.6540e6 | 2.6268e6 | 12101 | 116 | 1.03 | 118 | - | - | 172 | - |
7. Conclusions

We have benchmarked the CS-TSSOS hierarchy on a number of challenging AC-OPF cases and demonstrated that the 1.5th order relaxation is indeed useful in reducing the optimality gap and certifying global optimality of AC-OPF solutions. We hope that the computational results presented in this paper would convince people to think of CS-TSSOS as an alternative tool for certifying global optimality of solutions of large-scale POPs. One line of future research is to improve the efficiency of the CS-TSSOS relaxations by relying on more advanced chordal extension algorithms. We also plan to design suitable branch and bound algorithms to reach better accuracy results such as 0.1% or 0.01% global optimality for the AC-OPF problem. Eventually, such a current degree of optimality, i.e., 1% cannot be judged without considering uncertainty. For power systems of any size, the uncertainties of loads, generations, storage as well as configurations, impact upon steady-state operating points. In a further dedicated study, we shall extend our presented framework to handle uncertainties by means of sparsity-adapted versions of the existing approaches for robust/parametric polynomial optimization [17, 20].

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