A General Approach for Construction of Deterministic Compressive Sensing Matrices

Mohamad Mahdi Mohades, Mohamad Hossein Kahaei

1Department of Electronics and Communication Engineering, Iran University of Science and Technology, Tehran, Iran
E-mail: kahaei@iust.ac.ir

Abstract: In this paper, deterministic construction of measurement matrices in Compressive Sensing (CS) is considered. First, by employing the column replacement concept, a theorem for construction of large minimum distance linear codes containing all-one codewords is proposed. Then, by applying an existing theorem over these linear codes, deterministic sensing matrices are constructed. To evaluate this procedure, two examples of constructed sensing matrices are presented. The first example contains a matrix of size $p^2 \times p^2$ and coherence $1/p$, and the second one comprises a matrix with the size $p^{(p - 1)} \times p^{(p - 1)}$ and coherence $1/(p - 1)$, where $p$ is a prime integer. Based on the Welch bound, both examples asymptotically achieve optimal results. Moreover, by presenting a new theorem, the column replacement is used for resizing any sensing matrix to a greater-size sensing matrix whose coherence is calculated. Then, using an example, the outperformance of the proposed method is compared to a well-known method. Simulation results show the satisfying performance of the column replacement method either in created or resized sensing matrices.

1 Introduction

Compressive Sampling (CS) has widely been used to dramatically reduce the sampling frequency below the Nyquist rate [1]. This is effectively performed based on the sparse property of a signal vector $x_{m \times 1} \in \mathbb{R}^m$ whose sparsity order, $k$, indicates the number of its non-zero entries. Then, the CS problem is illustrated by $y_{m \times 1} = A_{m \times n} x_{n \times 1}$, where $y_{m \times 1} \in \mathbb{R}^m$ shows the measurement vector, $A_{m \times n}$ is the sampling matrix (or sensing matrix), and $m \ll n$. This sampling is worthwhile provided that the original signal $x_{n \times 1}$ can be uniquely and exactly reconstructed. The sparsest solution to this problem using $l_0$-norm minimization has been investigated in [2] which leads to an NP-hard problem [3]. To simplify this issue, an $l_1$-norm minimization solution has been proposed in [4] in which $A_{m \times n}$ must satisfy the so-called Restricted Isometry Property (RIP). Then, $x_{n \times 1}$ can be exactly and uniquely reconstructed by using $l_1$-norm minimization.

Definition 1. The matrix $A_{m \times n}$ satisfies the RIP of order $k$ with constant $\delta_k \in [0, 1)$, if for all $k$-sparse vectors $x_{n \times 1}$ we have the next inequality [4],

$$1 - \delta_k \leq \frac{\|A_{m \times n} x_{n \times 1}\|_2^2}{\|x_{n \times 1}\|_2^2} \leq 1 + \delta_k. \quad (1)$$

As mentioned, the signal $x_{n \times 1}$ sampled by the appropriate sampling matrix $A_{m \times n}$ could be reconstructed through solving the related $l_1$-norm minimization problem. The latter problem may also be solved using the greedy algorithms such as the Orthogonal Matching Pursuit (OMP) algorithm or its modifications [5]. At the beginning, Gaussian sensing matrices whose elements are independent and identically distributed (i.i.d.) were proposed which satisfy the RIP with a high probability [6]. For such matrices, the number of rows, columns, and the sparsity order are related to each other through the inequality $m \geq O(k \log n)$. On the other hand, deterministic sensing matrices may be used. Even though, due to the huge number of $k$-sparse signals, evaluation of the RIP for a given deterministic sensing matrix is frustrating, such a matrix can still satisfy the RIP if its coherence is low enough. A deterministic sensing matrix $A_{m \times n}$ of coherence $\mu_A$ satisfies the RIP of order $k$ with $\delta_k \leq \mu_A (k - 1)$ where $k \leq 1/\mu_A + 1$ [7]. The coherence of $A_{m \times n}$, whose columns are defined as $[a_{1}, \ldots, a_{N}]$, is given by

$$\mu_A = \max_{i \neq j} \frac{|\langle a_i, a_j \rangle|}{\|a_i\|_2 \|a_j\|_2} \quad 1 \leq i, j \leq N. \quad (2)$$

The minimum coherence of an arbitrary $A_{m \times n}$ is provided by the Welch bound expressed as [8]

$$\mu_A \geq \sqrt{\frac{n - m}{m \left(\frac{n - 1}{m}\right)}} \quad (3)$$

where $\mu_A \geq \sqrt{1/m}$ for $m \ll n$. This means that for deterministic sensing matrices, in terms of coherence, the RIP of sparsity $k = O(\sqrt{m})$ is satisfied which is known as the square root bottleneck. To construct deterministic sensing matrices with a low coherence, the DeVore’s construction was firstly introduced [9]. In this structure, binary sensing matrices of size $p^2 \times p^2$ and coherence of $r/p$ are used which satisfy the RIP of order $k < p/r + 1$ where $r$ shows the degree of a polynomial with coefficients in Galois field $\mathbb{F}_p$. By considering polynomials over finite projective spaces, instead of the finite field $\mathbb{F}_p$, extension of the DeVore’s construction has led to lower coherence sensing matrices [10]. The algebraic geometry codes [11] have also been utilized for deterministic sensing matrices construction. Other binary structures can be seen in [12], [13] and [14]. In [15], by combining two similar Reed-Solomon generator matrices through the tensor product and utilizing the generated codewords, a set of complex-valued deterministic sensing matrices with the size of $q^2 \times q^3$ and coherence $1/q$, where $q$ is any prime power, has been proposed. Using the Paley graphs, a set of deterministic sensing matrices of size $(q + 1)/2 \times q$ with coherence less than twice the Welch bound has been developed [16]. In [17], by employing BCH codes and creating large minimum distance codewords, some binary, bipolar and ternary deterministic sensing matrices are obtained. In [18], to overcome the limitation on matrices size, some methods are developed to resize the sensing matrices for complex valued sensing matrices. From the other aspect, we should note that the RIP is only a sufficient condition meaning that there might be some other matrices not satisfying this condition or even RIP-less, while can be used as sensing matrices [19].
In [20], some matrices which satisfy a weaker condition than the RIP, so-called statistical RIP, are proposed. In this paper, to construct low coherence sensing matrices, the theorem of employing large minimum distance linear codes containing all-one codeword is first introduced [18]. Then, we exhibit the procedure of employing the column replacement concept to create large minimum distance linear codes containing all-one codeword. Secondly, we present two examples in which the theorem of constructing sensing matrices is applied over our proposed matrices and we compute the coherence of the resulting matrices. Next, by comparing our results with some well-known sensing matrices, we exhibit the generality of the proposed method. These deterministic sensing matrices are of sizes $p^2 \times p^3$ and $p(p-1) \times p^3$ whose coherence are $1/p$ and $1/(p-1)$, respectively; where $p$ is a prime integer. The former example is quite similar to that of [15] for the prime integer $p$. About the second example, considering that of [18] matrices of size $(p^2 - 1) \times p^3$ and coherence $1/(p-1)$ are suggested, our proposed structure outweighs that of [18] in terms of coherence owing to the fact that our method considers the same coherence with lesser rows in matrix construction. It is notable that the mentioned examples are asymptotically optimal by the Welch bound. Finally, we utilize the column replacement method to resize the sensing matrices and calculate their coherence. By presenting an example, we show that for the same size of sensing matrices, our procedure results in better sensing matrices in the sense of coherence compared to the Kronecker product developed in [13]. The rest of paper is as follows. Section 2 contains the main results. In Section 3 using computer simulations, the performance of our sensing matrices is compared with some existing matrices. Section 4 concludes the paper.

2 Construction of Deterministic Sensing Matrices

To employ large minimum distance codes to generate low coherence deterministic sensing matrices, the following theorem applies.

Theorem 1. [18] Suppose that $\mathcal{C}[N, k, p, d_{\text{min}}]$ is a $p$-ary linear code over finite field $\mathbb{F}_p$ whose minimum distance is $d_{\text{min}}$ and all-one codeword belongs to the codewords' space. Now, assume that matrix $\hat{A}_{N \times p^{k-1}}$ is created by deploying each codeword as a column of this matrix in such a way that from any set of form $\{a, a \oplus 1, \ldots, a \oplus (p-1)\}$ exactly one codeword is selected, where $\oplus$ stands for element to element summation over finite field $\mathbb{F}_p$. Then, we create $\hat{A}_{N \times p^{k-1}}$ as

$$A = \frac{1}{\sqrt{N}} \left[ \epsilon_i^T \beta \hat{a}_{\alpha \beta} \right]_{i, \beta},$$

where $\hat{a}_{\alpha \beta}$ is an element of $\hat{A}_{N \times p^{k-1}}$. Then, the coherence of $A$ will be less than $\sqrt{p(p-1)/N-p^2d_{\text{min}}}/2N$.

This theorem enables us to think of low coherence sensing matrices provided that we can create large minimum distance matrix of codewords. We create such matrices based on column replacement of two matrices as follows.

Definition 2. [21] Consider the primary matrix $A \in \mathbb{F}^{r \times m}$, where $A=\{a_{ij}\}_{i,j}$, and the pattern matrix $P \in \{1, \ldots, m\}^{N \times n}$, where $P\{p_{ij}\}_{i,j}$. Then, the resulting column replacement of $A$ in $P$ is a matrix with the form of $C = B(P)a \{b_{ij} \} \in \mathbb{F}^{rN \times n}$, where $b_{ij}(\beta-1)_{r \times 1} \gamma_{1, \ldots, \gamma_r}$, for $1 \leq \beta \leq N$, $1 \leq \gamma \leq n$, and $1 \leq s \leq r$.

One can easily see that from [21] the following equation holds:

$$C = B(P)a = B\left(\{p_{ij}\}\right).$$

This means that column replacement is applied over each element of the pattern matrix. We use the column replacement to generate large minimum distance codes.

2.1 Generating Large Minimum Distance Codes

We apply the column replacement method over matrices of code to propose a new matrix. It is proved that the columns of this matrix are also codewords and their minimum distance is a function of the primary and pattern matrices. Therefore, under some criteria, the minimum distance can be large enough to apply Theorem 1. Using the column replacement method, we define and prove the following theorem to construct large minimum distance linear codes.

Theorem 2. Let $A_{N \times p^k}$ and $P_{N' \times p^{k'}}$ be matrices whose columns are codewords of the linear codes $\mathcal{C}[N, k, p, d_{\text{min}}]$ and $\mathcal{C}'[N', k', p', d'_{\text{min}}]$, respectively, where $d_{\text{min}}$ and $d'_{\text{min}}$ are the minimum distances of each defined linear code and $p$ is a prime integer. Moreover, suppose that both linear codes own all-one codeword. Then, the result of column replacement of $A_{N \times p^k}$ with $P_{N' \times p^{k'}}$ will be the matrix $\hat{B}(C_{N \times p^k} \cdot C_{N' \times p^{k'}})$ whose columns are codewords of a linear code containing all-one codeword with minimum distance of $N\cdot N' - \left((N' - d'_{\text{min}})N + d'_{\text{min}}(N - d_{\text{min}})\right)$.

Proof: We should note that it is natural to construct $\hat{B}(P_{N' \times p^{k'}})$ over finite field $\mathbb{F}_p$ to correspond each element of this matrix to a column of matrix $A_{N \times p^k}$. First, we prove that the columns of $C_{N \times p^k}$ are codewords. To do so, we must prove that any linear combination of every two columns will result in another column of this matrix. Thus, the columns belong to a linear subspace and are codewords. Suppose that $A_{N \times p^k}$, $P_{N' \times p^{k'}}$, and $C_{N \times p^k}$ have the following form

$$A_{N \times p^k} = \left[\begin{array}{c} a_{N \times 1}^{(k-1)} + \alpha(k-2) + \cdots + \alpha_1 + \alpha_0 \\vdots \\alpha_0 \end{array}\right]_{k \times 1},$$

$$P_{N' \times p^{k'}} = \left[\begin{array}{c} p_{N' \times 1}^{(k-1)} + \alpha(k-2) + \cdots + \alpha_1 + \alpha_0 \\vdots \\alpha_0 \end{array}\right]_{k \times 1},$$

$$C_{N \times p^k} = \left[\begin{array}{c} C_{N \times 1}^{(k-1)} + \alpha(k-2) + \cdots + \alpha_1 + \alpha_0 \\vdots \\alpha_0 \end{array}\right]_{k \times 1},$$

where

$$\alpha_{i} = a_{i1} \cdot \pi_{i1} + \cdots + a_{iN} \cdot \pi_{iN} + \epsilon_{i}$$

and $\pi_{ij}$ belongs to the finite field $\mathbb{F}_p$ with $1 \leq i \leq N'$ and $1 \leq j \leq p^{k'}$. It is obvious that the elements of $C_{N \times p^k}$ belong to $\mathbb{F}_p$ since the elements of $A_{N \times p^k}$ belong to this field. Moreover, in definition of $A_{N \times p^k}$ the superscript $i \pi_{ij} \cdot \epsilon_{i}$ is selected to corroborate the fact that the columns of $A_{N \times p^k}$ correspond to the elements of $\mathbb{F}_p$ whose elements are in the form of $i \pi_{ij} \cdot \epsilon_{i}$, and $\pi_{ij}$ belongs to the finite field $\mathbb{F}_p$ with $1 \leq i \leq N'$ and $1 \leq j \leq p^{k'}$. Therefore, such correspondence means that if an element of $P_{N' \times p^{k'}}$ is $\pi_{ij} \cdot \epsilon_{i}$, then the corresponding column of $A_{N \times p^k}$ is $i \pi_{ij} \cdot \epsilon_{i}$, and $\pi_{ij}$ belongs to the finite field $\mathbb{F}_p$ with $1 \leq i \leq N'$ and $1 \leq j \leq p^{k'}$. Such relationship will let us pursue our proof much easier. Now, consider an arbitrary combination of two arbitrary columns of $C_{N \times p^k}$ as

$$C'' = \beta_1 C_{N \times 1} + \beta_2 C_{N' \times 1},$$

where $\beta_1$ and $\beta_2$ are any element of $\mathbb{F}_p$ and $\text{c}_{N \times 1}$ and $\text{c}_{N' \times 1}$ are two columns of $C_{N \times p^k}$. Next, we should prove that $c''$ is

$$c'' = \beta_1 C_{N \times 1} + \beta_2 C_{N' \times 1},$$

where $\beta_1$ and $\beta_2$ are any element of $\mathbb{F}_p$ and $\text{c}_{N \times 1}$ and $\text{c}_{N' \times 1}$ are two columns of $C_{N \times p^k}$. Next, we should prove that $c''$ is...
also a column of $C_{N' \times p'k'}$. Suppose that by applying the column replacement over two columns of $P_{N' \times p'k'}$, namely, $p_{N' \times 1}$ and $p_{N' \times 1}$, $c_{N' \times 1}$ and $c_{N' \times 1}$ are obtained as

$$c_{N' \times 1}^m = B (p_{N' \times 1}) = B \begin{bmatrix} c_{k-1} & \cdots & c_0 \end{bmatrix}$$

(10)

where $B(.)$ presents the column replacement operator and

$$\begin{align*}
0 & \leq \varepsilon_{k-1}, \varepsilon_{k-2}, \ldots, \varepsilon_1, \varepsilon_0 \leq p - 1 \\
0 & \leq \tau_{k-1}, \tau_{k-2}, \ldots, \tau_1, \tau_0 \leq p - 1 .
\end{align*}$$

Using (10) and (11), we get

$$c'' = B (\beta_1 p_{N' \times 1} + \beta_2 p_{N' \times 1})$$

(12)

Considering the fact that $\beta_1 p_{N' \times 1} + \beta_2 p_{N' \times 1} \in P_{N' \times p'k'}$, it is inferred that $C_{N' \times p'k'}$ is a matrix whose columns are also codewords of a linear subspace. Note that although we used the characteristic of $\beta_1 p_{N' \times 1} + \beta_2 p_{N' \times 1}$ to find $d_{\text{min}}$, $\beta_1, \beta_2 \in F_p$, it can be applied to any $\beta_1, \beta_2$ in the given finite field $F_p$. To analytically calculate the minimum distance of codewords included in $C_{N' \times p'k'}$, we replace each column of $A_{N \times p}$ instead of each element of $P_{N' \times p'k'}$. In construction of each column of $C_{N' \times p'k'}$, the worst case happens when: 1) two columns of $P_{N' \times p'k'}$ have exactly $N' - d_{\text{min}}$ elements in common and 2) the rest of these elements are chosen in such a way that the corresponding columns of $A_{N \times p}$ have $N - d_{\text{min}}$ elements in common. By such an assumption, the resulting columns will have $(N' - d_{\text{min}}) + (N - d_{\text{min}})$ common elements which is the worst case, and therefore, the minimum distance of codewords of $C_{N' \times p'k'}$ is

$$d_{\text{min}} = N' - \left( (N' - d_{\text{min}}) + (N - d_{\text{min}}) \right)$$

(13)

The last step is to prove that $C_{N' \times p'k'}$ owns the all-one codeword. In the above theorem, it is assumed that $A_{N \times p}$ owns the all-one codeword which means that it is possible to consider the generator matrix $G_{N \times N}$ has a basis, say $g_{0N}$, whose all elements are one. Let the all-one column of $A_{N \times p}$ be

$$a_{N \times 1} = \begin{bmatrix} 0_0 \times (k-1) + 0_1 \times (k-3) + \cdots + 0_1 \times 1 \end{bmatrix} = g_0$$

(14)

Considering the fact that the elements of $P_{N' \times p'k'}$ belong to the finite field $F_p$, it can be shown that $P_{N' \times p'k'}$ has the all-one codeword, i.e.,

$$p_{N' \times 1} = \begin{bmatrix} 0_0 \times (k-1) + \cdots + 0_1 \times 1 \end{bmatrix} = \begin{bmatrix} 1 \end{bmatrix}$$

(15)

Using (14), (15), and the correspondence between the column vector $a_{N \times 1}$ and the element $0_0 \times (k-1) + \cdots + 0_1 \times 1$, one can immediately infer that $C_{N' \times p'k'}$ owns an all-one column as

$$B \left( p_{N' \times 1} \right) = \begin{bmatrix} 1 & \cdots & 1 \end{bmatrix} \in C_{N' \times 1}$$

(16)

This completes the proof.

In the following examples, using $C_{N' \times p'k'}$, we apply Theorem 2 to construct low coherence deterministic sensing matrices.

2.2 Case Study

To employ Theorems 1 and 2, we introduce some sensing matrices with low coherences and calculate their exact coherence value.

Example 1. Let the generator matrix of codewords of $A_{N \times p}$ and $P_{N' \times p'k'}$ be $G_{2 \times p} = \begin{bmatrix} g_0 \end{bmatrix}$, where $g_0 = [1, \ldots, 1]$ and $g_1 = [0, \ldots, p-1]$ are, Moreover, suppose that the columns of $A_{N \times p}$ and $P_{N' \times p'k'}$ are of the form $\gamma_0 g_0 + \gamma_1 g_1$, respectively, where $\gamma_0, \gamma_1 \in F_p$ and $\gamma_0, \gamma_1 \in F_p$. It is obvious that $A_{N \times p}$, $P_{N' \times p'k'}$, and $C_{N' \times p'k'}$ are of sizes $p \times p^2$, $p \times p^2$, and $p^2 \times p^4$, respectively. Trivially, the minimum distances of linear codes created by $G_{2 \times p}$ is $p - 1$ for both codewords $A_{p \times p}$ and $P_{p \times p}$; i.e., $d_{\text{min}} = p - 1$ and $d_{\text{min}} = p - 1$. The minimum distance of $C_{p^2 \times p^4}$ is thus equal to $d_{\text{min}} = (p - 1)^2$ which is a maximum minimum distance. Let us apply Theorem 1 to this matrix. We have the following equation:

$$C_{p^2 \times p^4} = \frac{1}{\sqrt{p^4}} \begin{bmatrix} c \end{bmatrix} \in \mathbb{R}_{p^t}$$

(17)

where $c_{t-t}$ is an element of $C_{p^2 \times p^4}$ and this matrix is constructed from matrix $C_{p^2 \times p^4}$ such that from each column set $c_{p^2 \times 1} \in C_{p^2 \times 1} = \mathbb{R}_{p^2}$ only one column is selected. We allow that the coherence of $C_{p^2 \times p^4}$ is $1/p$.

Proof: First, we can write,

$$C_{p^2 \times p^4} = B \left( P_{p \times p^4} \right)$$

(18)

As illustrated in (16), the all-one column of $C_{p^2 \times p^4}$ is obtained when the column replacement acts on the all-one column of $P_{p \times p^4}$ as

$$B \left( p_{p \times 1} \right) = \begin{bmatrix} 1 & \cdots & 1 \end{bmatrix} \in \mathbb{F}_{p^2}$$

(19)

Imagine the column $c_{p^2 \times 1}$ of $C_{p^2 \times p^4}$ has the following form:

$$c_{p^2 \times 1} = B \left( \left( \lambda_1 \alpha + \lambda_0 \right) g_1 + \left( \lambda_0 \alpha + \lambda_0 \right) g_0 \right)$$

(20)

where $\lambda_1, \lambda_0, \lambda_0, \lambda_0 \in F_p$. Using the definition of $C_{p^2 \times p^4}$, $c_{p^2 \times 1}$, and Theorem 2 we can write

$$c_{p^2 \times 1} = \beta_0 (1) p_{p \times 1} = B \left( \left( \lambda_1 \alpha + \lambda_0 \right) g_1 + \left( \lambda_0 \alpha + \lambda_0 \right) g_0 \right)$$

(21)

Obviously to choose one column of the column set

$$\{ c_{p^2 \times 1}, c_{p^2 \times 1} + 1_{p \times 1}, \ldots, c_{p^2 \times 1} + (p - 1)_{p \times 1} \}$$

it is enough to keep $\lambda_0$ constant while applying the column replacement over $P_{p \times p^4}$. For convenience, we consider $\lambda_0 = 0$. Then, the relationship between $C_{p^2 \times p^4}$ and $P_{p \times p^4}$ is given by

$$C_{p^2 \times p^4} = B \left( P_{p \times p^4} \right)$$

(22)

where $P_{p \times p^4}$ contains those columns of $P_{p \times p^4}$ which correspond to $(\lambda_1 \alpha + \lambda_0) g_1 + (\lambda_0 \alpha) g_0$. So far, we could show how
it is feasible to generate $C_p' p^2 \times p^3$ based on $P_p p^2 p'$. Let us proceed to construct the sensing matrix $C'' p^2 \times p^3$ based on $[17]$ as

$$[f (c''_t)]_{r,t} = \begin{bmatrix} c''_{p^2 \times 1} & \ldots & c''_{p^2 \times 1} \end{bmatrix}$$

(23)

where $f (c''_t) = \frac{1}{\sqrt{p'}} e^{j2\pi \phi t}$ is defined for simplicity. Since the elements of $C'' p^2 p^3$ come from the primary matrix $A_{p^2 p^2}$, i.e. $c''_t \in A_p p^2 p^2$, we can firstly apply the transformation function $f$ on each element of $A_{p^2 p^2}$ which leads to

$$\tilde{A}_{p^2 p^2} = \begin{bmatrix} \tilde{a}_{p^2 \times 1}^0 + 0 \cdot \tilde{a}_{p^2 \times 1}^1, \ldots, \tilde{a}_{p^2 \times 1}^\gamma_1 + \gamma_0 \tilde{a}_{p^2 \times 1}^1, \ldots, \tilde{a}_{p^2 \times 1}^{p-1} \alpha + (p-1) \tilde{a}_{p^2 \times 1}^1 \end{bmatrix}$$

(24)

where $\tilde{a}_{p^2 \times 1}^\gamma_1 + \gamma_0$ is any arbitrary column of $\tilde{A}_{p^2 p^2}$ which corresponds to the column $a_{p^2 \times 1}^\gamma_1 + \gamma_0$ of matrix $A_{p^2 p^2}$, and superscript $\gamma_1 + \gamma_0$ is assumed to accommodate the correspondence of columns of $\tilde{A}_{p^2 p^2}$ and the elements of $P_p p^2 p'$. To simplify the notations, we have considered the defined transformation function $f$ on matrices and vectors as

$$\tilde{A}_{p^2 p^2} = f (A_{p^2 p^2})$$

(25)

Then, we can rewrite $C'' p^2 p^3$ as

$$C'' p^2 p^3 = B' (P_p p^3)$$

(26)

where $B'$ represents the column replacement for $\tilde{A}_{p^2 p^2}$. Now, we can calculate the coherence of $C'' p^2 p^3$ whose coherence is

$$\mu_{C''} = \max_{j \neq k} \left| \left( c''_{p^2 \times 1}, c''_{p^2 \times 1} \right) \right|$$

(27)

where $(\cdot, \cdot)$ illustrates the inner product of two vectors and $c''_{p^2 \times 1}$, $c''_{p^2 \times 1}$ are two different columns of $C'' p^2 p^3$. Considering $c''_{p^2 \times 1}$, $c''_{p^2 \times 1}$ as

$$c''_{p^2 \times 1} = B' \left( P_p p^3 \right)$$

(28)

where $P_p p^3$ and $P_p p^3$ are the $j^{th}$ and $k^{th}$ columns of $P_p p^2 p'$, respectively. By rewriting the corresponding columns of $P_p p^2 p'$, $\mu_{C''}$ reduces to

$$\mu_{C''} = \max_{j \neq k} \left| \begin{bmatrix} (\lambda_1, \lambda_0), \lambda_1^{p-1} \lambda_0, \lambda_1 \end{bmatrix} \right|$$

(29)

Owing to the fact that coherence is considered over different columns, $\lambda_1, \lambda_0, \lambda_1^{p-1} \lambda_0$ cannot be zero at the same time. Considering $\lambda_1 = \lambda_1^{p-1} - \lambda_1$, $\lambda_0 = \lambda_0^{p-1} - \lambda_1$, and $\lambda_1^{p-1}, \lambda_0^{p-1}, \lambda_1, \lambda_0 \in \mathbb{R}$. Therefore, $\mu_{C''}$ can be stated as

$$\mu_{C''} = \max_{\lambda_1, \lambda_0, \lambda_1^{p-1} \lambda_0, \lambda_1^{p-1} \lambda_0} \left| \begin{bmatrix} f (\lambda_1, \lambda_0), \lambda_1, \lambda_0, \lambda_1^{p-1} \lambda_0, \lambda_1^{p-1} \lambda_0 \end{bmatrix} \right|$$

(30)

where $\phi \in \mathbb{R}$. Using $f$, $\mu_{C''}$ becomes

$$\mu_{C''} = \max_{\lambda_1, \lambda_0, \lambda_1^{p-1} \lambda_0, \lambda_1^{p-1} \lambda_0} \left| \begin{bmatrix} f (\lambda_1, \lambda_0), \lambda_1, \lambda_0, \lambda_1^{p-1} \lambda_0, \lambda_1^{p-1} \lambda_0 \end{bmatrix} \right|$$

(31)

where $\ast$ indicates complex conjugation. Then, the inner product of $f$ can be simplified as

$$\left< f (u_1), f^* (u_2) \right> = \left< f (u_1) - u_2 \right>$$

(32)

Moreover, according to $[32]$, the inner term of transformation function $f$ in $[32]$ for an arbitrary $\phi \in \mathbb{R}$ can be rewritten as

$$\left< f (\lambda_1, \lambda_0), \lambda_1, \lambda_0, \lambda_1^{p-1} \lambda_0, \lambda_1^{p-1} \lambda_0 \right> = \left< f (\lambda_1, \lambda_0), \lambda_1, \lambda_0, \lambda_1^{p-1} \lambda_0, \lambda_1^{p-1} \lambda_0 \right>$$

(33)

where $\lambda_1 = \lambda_1^{p-1} - \lambda_1$, $\lambda_0 = \lambda_0^{p-1} - \lambda_1$, and $\lambda_1^{p-1}, \lambda_0^{p-1}, \lambda_1, \lambda_0 \in \mathbb{R}$. Therefore, $\mu_{C''}$ can be stated as

$$\mu_{C''} = \max_{\lambda_1, \lambda_0, \lambda_1^{p-1} \lambda_0, \lambda_1^{p-1} \lambda_0} \left| \begin{bmatrix} f (\lambda_1, \lambda_0), \lambda_1, \lambda_0, \lambda_1^{p-1} \lambda_0, \lambda_1^{p-1} \lambda_0 \end{bmatrix} \right|$$

(34)
\[\mu_{C\nu} = \max_{\lambda''_{p,\nu} \in \mathbb{F}_p} \frac{1}{\nu} \sum_{\phi \neq 0} f(0, \lambda''_{p,\nu}) + \cdots + f(\gamma \lambda''_{p,\nu}) + \cdots + f((p-1)\lambda''_{p,\nu}) + \cdots + f(\gamma \lambda''_{1,\nu}) + \cdots + f((p-1)\lambda''_{1,\nu})\]

\[f(\lambda''_{p,\nu}) + \cdots + f((\lambda''_{1,\nu} + \lambda''_{0,\nu}) \gamma + \lambda''_{0,\nu}) + \cdots + f((\lambda''_{1,\nu} + \lambda''_{0,\nu}) (p-1) + \lambda''_{0,\nu})\]

\[= \max_{\lambda''_{p,\nu} \in \mathbb{F}_p} \frac{1}{\nu} \sum_{\phi = 0, \gamma = 0} f \left( \gamma^2 \lambda''_{p,\nu} \right) \text{ for } \phi = \nu^{\phi} \]

Example 2. Considering the same procedure as Example 1 we only highlight the changes. Let the generator matrix of the codewords of \(A_{N \times p^2}\) and \(P_{N \times p^2}^T\) be \(G_{2 \times p} = [g_0, g_1]^T\) and \(G'_{2 \times p} = \left[ g_{0}', g_{1}' \right]^T\), respectively, where \(g_0 = [1, \ldots, 1]^T_{1 \times p}\), \(g_1 = [0, \ldots, p-1]^T_{1 \times (p-1)}\), \(g_{0}' = [1, \ldots, 1]^T_{1 \times (p-1)}\), and \(g_{1}' = [0, \ldots, p-2]^T_{1 \times (p-1)}\). The canonical additive character of finite field \(F_{p^2}\) is \(\chi(\gamma) = e^{2\pi i p \gamma}, \quad \forall \gamma \in \mathbb{F}_p\) (36)

Theorem 3. Let \(\{a_k\}_{k=1}^n\) and \(\{b_k\}_{k=1}^n\) be two arbitrary matrices in \(\mathbb{F}_{p^2}\) and \(\mathbb{F}_p\), respectively.]

\[X = Z \oplus Y = \begin{bmatrix} z_{11} Y & \cdots & z_{1n} Y \\ \vdots & \ddots & \vdots \\ z_{m1} Y & \cdots & z_{mn} Y \end{bmatrix} \quad (40)\]

where \(Z = \begin{bmatrix} z_{11} & \cdots & z_{1n} \\ \vdots & \ddots & \vdots \\ z_{m1} & \cdots & z_{mn} \end{bmatrix}\) and \(Y\) are two arbitrary matrices.

2.3 Matrix Resizing

Here, we employ the column replacement concept to resize any deterministic sensing matrix and illustrate that after resizing to a matrix with a given greater size, our method leads to a matrix with lower coherence than that of [18]. To do so, first we indicate Theorem 3 from [18] and then develop our method in Theorem 4.

Theorem 4. Let \(A_{N \times N} = \begin{bmatrix} a_{m1} & \cdots & a_{mn} \end{bmatrix}\) be a deterministic sensing matrix with coherence of \(\mu_A\), where superscript \(t\) is to create a correspondence between the primary and pattern matrices in column replacement and \(a_{i \times 1}\) are normalized vectors for \(1 \leq i \leq N\). Furthermore, let \(P_{N \times L} = \left[ p_{01} \cdots p_{L1} \right] \) be a matrix whose elements are \(t_{ij}\), \(1 \leq j \leq n\) and no two distinct columns of this matrix have more than \(d_{p} \) elements in common. Trivially, \(1 \leq i \leq L\) and \(0 \leq d_{p} \leq N-1\) is true statements. Then, the coherence of \(C_{MN \times L} = \frac{1}{\sqrt{N}} B(P_{N \times L})\) is \(\mu_C \leq \frac{1}{\sqrt{N}} (d_{p} + (N-d_{p}) \mu_A)\), where \(B(P_{N \times L}) \) means column replacement of \(A_{N \times N}\) based on the pattern matrix \(P_{N \times L}\), and the factor \(\frac{1}{\sqrt{N}}\) normalizes the columns after column replacement.

Proof: Suppose that \(C_{MN \times L} = \begin{bmatrix} c_{m1} & \cdots & c_{mL} \end{bmatrix}\). Then, \(\mu_C\) is defined as

\[\mu_C = \max_{1 \leq i \leq L, \mu \leq L} \left| \begin{bmatrix} c_{i1}^1 \cdots c_{i1}^L \\ \vdots \\ c_{in}^1 \cdots c_{in}^L \end{bmatrix} \right| \quad (41)\]

Assume that the two arbitrary columns \(c_{m1}^1 \cdots c_{m1}^L\) and \(c_{m2}^1 \cdots c_{m2}^L\) are constructed by employing column replacement over two columns \(p_{1}^i \cdots p_{L1}^i \) of matrix \(P_{N \times L}\), as

\[c_{m1}^i = \frac{1}{\sqrt{N}} \left[ B\left( p_{i1}^1 \right) \cdots B\left( p_{i1}^L \right) \right] \quad (42)\]

where \(1 \leq s \leq N\) and \(1 \leq i_1, i_2 \leq L\). Then, we have...
In common; i.e, $d_\mathbf{P} = 1$. Consequently, the sensing matrix $\mathbf{C}$ which can be found from Theorem 3 based on $A_{p^3 \times p^3}$ and $\mathbf{P}_{p^3 \times p^3}$ is a matrix of size $p^2 \times p^3$ and coherence $\mu_C = (2p-1)p^2$. For comparison purposes, we assume that the size of Kronecker product matrix $\mathbf{C}^*$ is $p^2 \times p^3$. According to Theorem 3, for $\mathbf{A}_{p^3 \times p^3}$ we need a matrix of size $p \times p^3$ to reach the same size of $\mathbf{C}$. In other words, $\mathbf{B}$ is a sensing matrix of size $p \times p^3$ whose coherence according to the Welch bound cannot be less than $\mu_{\mathbf{B}} = \sqrt{\frac{p+1}{p^3} + \frac{1}{p^2}}$. Therefore,

$$\mu_{\mathbf{C}^*} = \max (\mu_{\mathbf{A}}, \mu_{\mathbf{B}}) = \max \left( \frac{1}{p^2} \sqrt{\frac{p+1}{p^2} + \frac{1}{p^3}}, \mu_{\mathbf{C}} \right)$$

Consider the following comparison,

$$\sqrt{\frac{p+1}{p^2} + \frac{1}{p^3}} \geq \frac{1}{p} \Rightarrow p^2 \frac{p+1}{p^3} \geq p^2 + p^3 + 1$$

Obviously, $p^3 > p + 1$ for any prime integer, and therefore,

$$\mu_{\mathbf{C}^*} = \sqrt{\frac{p+1}{p^2} + \frac{1}{p^3}}$$

Now, we compare the coherence of the proposed sensing matrix obtained through resizing by the column replacement method and that of [18]. Therefore, comparison between $\mu_{\mathbf{C}}$ and $\mu_{\mathbf{C}^*}$ results in

$$(2p-1)p^2 \geq \sqrt{\frac{p+1}{p^2} + \frac{1}{p^3}} \Rightarrow 0 < \left( p^4 - p \right) (p - 3) + 1$$

From [54], one can easily see that for the prime integers more than 3, the proposed method results in a lower coherence. Note that we assumed $\mathbf{B}$ is a sensing matrix of size $p \times p^3$ attaining the Welch bound. However, to the best of our knowledge, the best deterministic sensing matrices asymptotically attaining the Welch bound are of size $p \times p^2$. This result again corroborates our method as a better approach compared to the Kronecker product method in [18].

### 3 simulation results

We compare the proposed method with those of [18] and random sensing matrices. To do so, we consider three different scenarios to verify all the examples presented in Section III. The results are shown based on the average of 3000 independent trials (1000 trials for Figs. 3 and 4) for different $k$-sparse signals. In each scenario, the measurement matrix performance is verified based on the recovery percentage as well as the output SNR, whereas the sparsity order is increasing and the input SNR is augmented from 0 to 100 dB. The reconstruction algorithm is Orthogonal Matching Pursuit which is suitable to solve $l_1$-minimization problems. In the first problem, we examine Example 1 in which matrices of size $p^2 \times p^3$ and coherence $1/p$ are generated over finite field $\mathbb{F}_p$. We get a sensing matrix of size $25 \times 125$ and coherence of $1/5$. Over finite field $\mathbb{F}_p$, the construction of [18] leads to a size of $24 \times 125$ with coherence of $1/4$. The third matrix in this scenario is a Gaussian sensing matrix of size $25 \times 125$. The results are depicted in Figs. 1 and 2. Sparsity order of 12 has been considered in Fig. 2. Apparently, our proposed deterministic sensing matrices act better than that of [18]. Meanwhile, both deterministic structures surpass random sensing matrices in the sense of recovery percentage as well as the output SNR.

In the last scenario, we compare our method to resize existing sensing matrices with the Kronecker product method. We assume that the primary sensing matrix $\mathbf{A}$ is of size $25 \times 125$ and coherence $1/5$ and according to our example we need a pattern matrix of size $5 \times 5 \times 125$ for each two distinct columns of which there exists at most one element in common and the resulting sensing matrix will be of size $25 \times 15625$ and coherence 0.36. On the other hand, to obtain a similar size sensing matrix based on the Kronecker product method, we need a matrix of size $5 \times 125$ with near optimal
coherence. Such a matrix has not been introduced so far and the best known structure attaining the Welch bound with the same number of rows is a matrix of size $5 \times 25$ and coherence of $1/\sqrt{5}$. Therefore, we resize $A$ up to $125 \times 3125$ for which the resulting matrix has a coherence of $1/\sqrt{5}$. Although, it is severely unfair to our construction; in which we could obtain a matrix of size $125 \times 15625$ and coherence $0.36$, for comparison with the Kronecker product method, we have to reduce its size to $125 \times 3125$. We also consider a Gaussian sensing matrix of size $125 \times 3125$. Finally, Figs. 3 and 4, where the sparsity is fixed to 28 to generate Fig. 5, illustrate the results of the third scenario in which we proposed a new method for resizing deterministic sensing matrices. As seen, in this example our method performs far better than the Kronecker product method. Notice that we have only kept $1/5$ of the available columns in our structure which reveals another advantage of the proposed method.

4 Conclusion

The column replacement concept was used to construct low coherence deterministic sensing matrices for CS problems. First, this method was employed to create deterministic sensing matrices. It was applied over two matrices of linear codewords to generate a new matrix of linear codewords suitable for creating sensing matrices. By using two examples and comparing with some existing asymptotically optimal matrices, we illustrated the generality of the proposed method. Secondly, we exploited the column replacement method to resize existing sensing matrices and found a bound for the coherence of resulted matrices. Using an example, as an especial case, it was shown that this method behaves better than the Kronecker product method developed in [18]. Observe that only a special case of Theorem 2 results in the matrices proposed in [15], and hence, one may find much more diverse and applicable structures by manipulating Theorem 2. Moreover, as seen in Example 3 our resizing approach is not only a way to increase the size of sampling matrices but also an approach to obtain well-structured sensing matrices. Finally, simulation results showed that by utilizing the proposed column replacement method, the constructed matrices perform better than the existing deterministic sensing matrices with approximately the same size. Moreover, simulation results confirmed that the suggested procedure of resizing sensing matrices can perform much better compared to the Kronecker product method.

APPENDIX A

Consider the coherence as follows.

$$
\mu_{C''} = \max_{\lambda''_{01}, \lambda''_{12}, \lambda''_{11} \in \mathbb{F}_p} \frac{1}{p^2} \left| \sum_{\gamma=0}^{p-1} \sum_{\phi=0}^{p-1} e^{j2\pi \gamma (\lambda''_{11} \phi + \lambda''_{01})} \right|
$$

Contemplating (A-1), the inner summation of (A-1) is non-zero, iff:

$$
\lambda''_{11} \phi + \lambda''_{01} = 0
$$

Consider the coherence as follows.

$$
\mu_{C''} = \max_{\lambda''_{01}, \lambda''_{12}, \lambda''_{11} \in \mathbb{F}_p} \frac{1}{p^2} \left| \sum_{\gamma=0}^{p-1} \sum_{\phi=0}^{p-1} e^{j2\pi \gamma (\lambda''_{11} \phi + \lambda''_{01})} \right|
$$

Contemplating (A-1), the inner summation of (A-1) is non-zero, iff:

$$
\lambda''_{11} \phi + \lambda''_{01} = 0
$$

APPENDIX A

Consider the coherence as follows.
Then, the following cases will happen:

1. \((\lambda''_{11}, \lambda''_{01}) = (0, 0)\) and \(\lambda_{10}'' \neq 0\), for which \((\ref{eq:lambda''})\) reduces to

\[
\mu_{C''} = \max_{(\lambda''_{10} \neq 0) \in \mathbb{F}_p} \frac{1}{p} \left| \sum_{\phi = 0}^{p-1} e^{2\pi i \phi \lambda''_{10}} \right| = 0 \tag{A-3}
\]

In this case, due to the non-zero value of \(\lambda''_{10}\), coherence of \(\phi''\) is zero.

2. \(\lambda''_{11} \neq 0\) and \(\lambda_{10}'' \neq 0\) are arbitrarily chosen. Consequently, \(\lambda''_{11} \phi + \lambda''_{01}\) is zero only for one value of \(\phi\), say \(\phi = \phi_1\). Therefore, the coherence value is given by

\[
\mu_{C''} = \max_{\phi \neq \phi_1, \lambda''_{10} \in \mathbb{F}_p} \frac{1}{p} \left| \sum_{\phi = 0}^{p-1} \sum_{\gamma = 0}^{p-1} e^{2\pi i \phi \lambda''_{01}} \chi(\gamma(\lambda''_{11} \phi + \lambda''_{01} + \lambda_{10}'')) \right| = 0 \tag{B-4}
\]

Based on the above cases, one can infer that the coherence of \(C''\) is \(1/(p-1)\), i.e. \(\mu_{C''} = 1/(p-1)\).

\section{References}

1. D. L. Donoho, “Compressed sensing,” Information Theory, IEEE Transactions on, vol. 52, no. 4, pp. 1289–1306, 2006.
2. E. J. Candès and T. Tao, “Near-optimal signal recovery from random projections: Universal encoding strategies?,” Information Theory, IEEE Transactions on, vol. 52, no. 12, pp. 5443–5456, 2006.
3. D. Ge, X. Jiang, and Y. Ye, “A note on the complexity of l p minimization,” Mathematical programming, vol. 129, no. 2, pp. 285–299, 2011.
4. E. J. Candès and T. Tao, “Decoding by linear programming,” Information Theory, IEEE Transactions on, vol. 51, no. 12, pp. 4203–4215, 2005.
5. J. A. Tropp and A. C. Gilbert, “Signal recovery from random measurements via orthogonal matching pursuit,” Information Theory, IEEE Transactions on, vol. 53, no. 12, pp. 4655–4666, 2007.
6. R. Bariunik, J. Moufan, R. Devore, and M. Wakin, “A simple proof of the restricted isometry property for random matrices,” Constructive Approximation, vol. 26, no. 3, pp. 253–263, 2008.
7. J. Bourgain, S. Dilworth, K. Ford, S. Konyagin, D. Kutzarova, et al., “Explicit constructions of rip matrices and related problems,” Duke Mathematical Journal, vol. 159, no. 1, pp. 145–185, 2011.
8. L. R. Welch, “Lower bounds on the maximum cross correlation of signals (con- corr.),” Information Theory, IEEE Transactions on, vol. 20, no. 3, pp. 397–399, 1974.
9. R. A. DeVore, “Deterministic constructions of compressed sensing matrices,” Journal of Complexity, vol. 23, no. 4, pp. 918–925, 2007.
10. A. Mohades and A. Tadaion, “Finite projective spaces in deterministic construction of measurement matrices,” IET Signal Processing, 2016.
11. S. Li, F. Guo, G. Ge, and S. Zhang, “Deterministic construction of compressed sensing matrices via algebraic curves,” Information Theory, IEEE Transactions on, vol. 58, no. 8, pp. 5035–5044, 2012.
12. P. Sasmal, R. R. Naidu, C. S. Sastry, and P. Jampana, “Compilation of binary compressed sensing matrices,” IEEE Signal Processing Letters, vol. 23, no. 8, pp. 1096–1100, 2016.
13. R. R. Naidu, P. Jampana, and C. S. Sastry, “Deterministic compressed sensing matrices: Construction via euler squares and applications,” IEEE Trans. Signal Processing, vol. 64, no. 14, pp. 3566–3575, 2016.
14. R. R. Naidu and C. R. Murthy, “Construction of binary sensing matrices using extremal set theory,” IEEE Signal Processing Letters, vol. 24, no. 2, pp. 211–215, 2017.
15. M. Mohades, A. Mohades, and A. Tadaion, “A red-soledon soluble code based measurement matrix with small coherence,” Signal Processing Letters, IEEE, vol. 21, no. 7, pp. 839–843, 2014.
16 A. Amini, H. Bagh-Sheikhi, and F. Marvasti, “From paley graphs to deterministic sensing matrices with real-valued gramians,” in Sampling Theory and Applications (SampTA), 2015 International Conference on, pp. 372–376, IEEE, 2015.

17 A. Amini and F. Marvasti, “Deterministic construction of binary, bipolar, and ternary compressed sensing matrices,” Information Theory, IEEE Transactions on, vol. 57, no. 4, pp. 2360–2370, 2011.

18 A. Amini, V. Montazerhodjat, and F. Marvasti, “Matrices with small coherence using-ary block codes,” Signal Processing, IEEE Transactions on, vol. 60, no. 1, pp. 172–181, 2012.

19 W. Xu and B. Hassibi, “Efficient compressive sensing with deterministic guarantees using expander graphs,” in Information Theory Workshop, 2007. ITW’07. IEEE, pp. 414–419, IEEE, 2007.

20 R. Calderbank, S. Howard, and S. Jafarpour, “Construction of a large class of deterministic sensing matrices that satisfy a statistical isometry property,” Selected Topics in Signal Processing, IEEE Journal of, vol. 4, no. 2, pp. 358–374, 2010.

21 C. J. Colbourn, D. Horsley, and C. McLean, “Compressive sensing matrices and hash families,” Communications, IEEE Transactions on, vol. 59, no. 7, pp. 1840–1845, 2011.

22 R. Lidl and H. Niederreiter, Finite fields, (Cambridge university press, Cambridge, 2nd edn. 1997).