Some Results on the Vector Gaussian Hypothesis Testing Problem

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Abstract—This paper studies the problem of discriminating two multivariate Gaussian distributions in a distributed manner. Specifically, it characterizes in a special case the optimal type-II error exponent as a function of the available communication rate. As a side-result, the paper also presents the optimal type-II error exponent of a slight generalization of the hypothesis testing against conditional independence problem where the marginal distributions under the two hypotheses can be different.

I. INTRODUCTION

Consider the single-sensor single-detector hypothesis testing scenario in Fig. 1. The sensor observes a source sequence $X^n \triangleq (X_1, \ldots, X_n)$ and communicates with the detector, who observes source sequence $Y^n \triangleq (Y_1, \ldots, Y_n)$, over a noise-free bit-pipe of rate $R \geq 0$. Here, $n$ is a positive integer that denotes the blocklength and the sequence of pairs $\{(X_t, Y_t)\}_{t=1}^n$ is independent and identically distributed (i.i.d) according to a jointly Gaussian distribution of zero-mean and of joint covariance matrix that depends on the hypothesis $H \in \{0, 1\}$. Under hypothesis

$H = 0 : \left\{ \left( X_t, Y_t \right) \right\}_{t=1}^n \text{ i.i.d. } \sim \mathcal{N}(0, K), \tag{1}$

and under hypothesis

$H = 1 : \left\{ \left( X_t, Y_t \right) \right\}_{t=1}^n \text{ i.i.d. } \sim \mathcal{N}(0, K). \tag{2}$

Based on its observations $Y^n$ and the message it receives from the sensor, the Detector decides on the hypothesis by producing $\hat{H} \in \{0, 1\}$. The goal of this decision is to maximize the exponential decrease of the probability of type-II error (i.e., guessing $\hat{H} = 0$ when $H = 1$), while ensuring that the probability of type-I error (i.e., guessing $\hat{H} = 1$ when $H = 0$) goes to zero as $n \to \infty$.

The described single-sensor single-detector problem has previously been studied in [1]–[4] for various joint distributions on the i.i.d. observations. In particular, [4] identified the largest type-II exponent that is achievable in a setup that they termed testing against conditional independence. An explicit expression for the vector Gaussian case was recently found in [5] (see Theorem 2 therein which actually provides the solution of a more general, distributed, setting). For all other cases a computable single-letter characterization of the largest achievable type-II error exponent remains open. This line of works has also been extended to multiple sensors [6], multiple detectors [7], [8], interactive terminals [9]–[11], multi-hop networks [12]–[16], noisy channels [17], [19] and to scenarios with privacy constraints [19]–[21].

In this paper we present a computable single-letter characterization of the largest type-II error exponent achievable for the Gaussian vector hypothesis testing problem for a class of matrices $K$ and $K$. Our converse proof starts from the known multi-letter expression for this problem [1] and connects it to related results. The achievability proof is based on the coding scheme proposed in [3].

We end this introductory section with some remarks on notation. When two random variables $(X, Y)$ are independent given a third random variable $Z$ (i.e. $P_{XY|Z} = P_{X|Z}P_{Y|Z}$), we say $(X, Z, Y)$ form a Markov chain and write $X \leftrightarrow Z \leftrightarrow Y$. Both $D(P_X||P_{XY})$ and $D(X||X)$ denote the Kullback-Leibler divergence between two pmfs $P_X$ and $P_{XY}$. $h(\cdot)$, $I(\cdot;\cdot)$ and $I(\cdot;\cdot;\cdot)$ denote continuous entropy, mutual information and conditional mutual information. The set of all real numbers is denoted by $\mathbb{R}$. Boldface upper case letters denote random vectors or deterministic matrices, e.g., $X$, where the context should make the distinction clear. We denote the covariance matrix of a real-valued vector $X$ with distribution $P_X$ by $K_X = \mathbb{E}_{P_X}[XX^\dagger]$, where $\dagger$ indicates the transpose operation. Similarly, we denote the cross-correlation of two zero-mean vectors $X$ and $Y$ with joint distribution $P_{XY}$ by $K_{XY} = \mathbb{E}_{P_{XY}}[XY^\dagger]$, the conditional covariance matrix of $X$ given $Y$ with p.d.f $P_{XY}$ and with p.d.f $P_{X|Y}$ by $K_{X|Y} = \mathbb{E}_{P_{XY}}[XX^\dagger|Y]$ and $K_{X|Y} = \mathbb{E}_{P_{XY}}[XX^\dagger|Y]$, respectively. Finally, for a matrix $M$, we denote its inverse (if it exists) by $M^{-1}$ its determinant (if it exists) by $|M|$, its Moore-Penrose pseudo-inverse by $M^+$ and its pseudo-determinant by $|M|^+$.  

II. FORMAL PROBLEM STATEMENT

The sequences $X^n$ and $Y^n$ are as described before, where

\[ X^n \quad \text{Sensor} \quad M \in \{1, \ldots, W_n\} \quad \text{Detector} \quad \hat{H} \in \{0, 1\} \]

Fig. 1. Vector Gaussian hypothesis testing problem

we denote by $m$ the dimension of each vector $X_t$ and by $q$
the dimension of each vector $Y_t$. The Sensor, which observes $X^n$ applies an encoding function
\[ \phi_n: \mathbb{R}^{m \times n} \rightarrow M = \{1, \ldots, W_n\} \tag{3} \]
to this sequence and sends the resulting index
\[ M = \phi_n(X^n) \tag{4} \]
to the detector. Based on this message $M$ and its observation $Y^n$, the detector then applies a decision function
\[ \psi_n: M \times \mathbb{R}^{q \times n} \rightarrow \{0, 1\} \tag{5} \]
to decide on the hypothesis
\[ \hat{H} = \psi_n(M,Y^n). \tag{6} \]
The Type-I and type-II error probabilities at the detector are defined as:
\[ \alpha_n \triangleq \Pr\{\hat{H} = 1 | H = 0\} \tag{7} \]
\[ \beta_n \triangleq \Pr\{\hat{H} = 0 | H = 1\}. \tag{8} \]

**Definition 1.** Given rate $R \geq 0$, an error-exponent $\theta$ is said achievable if for all blocklengths $n$ there exist functions $\phi_n$ and $\psi_n$ as in (3) and (5) so that the following limits hold:
\[ \lim_{n \to \infty} \alpha_n = 0, \tag{9a} \]
\[ \theta \leq \lim_{n \to \infty} -\frac{1}{n} \log \beta_n \tag{9b} \]
and
\[ \lim_{n \to \infty} -\frac{1}{n} \log_2 W_n \leq R. \tag{9c} \]

**Definition 2** (Exponent-rate function). For any rate $R \geq 0$, the exponent-rate function $E(R)$ is the supremum of the set of all achievable error-exponents.

In essence, the problem of vector Gaussian hypothesis testing that we study here amounts to discriminating two covariance matrices. As we already mentioned the solution of this problem is known only in few special cases, namely the cases of testing against independence and testing against conditional independence [1], [4], and [5, Theorem 2].

### III. Optimal exponent for a class of covariance matrices

Let $K_X$ and $K_Y$ be $m$-by-$m$ dimensional matrices, $K_Y$ and $K_Y$ be $q$-by-$q$ dimensional matrices, and $K_{XY}$ and $K_{XY}$ be $m$-by-$q$ dimensional matrices such that
\[ K = \begin{bmatrix} K_X & K_{XY} \\ K_{XY}^\dagger & K_Y \end{bmatrix} \quad \text{and} \quad \bar{K} = \begin{bmatrix} \bar{K}_X & \bar{K}_{XY} \\ \bar{K}_{XY}^\dagger & \bar{K}_Y \end{bmatrix}. \tag{10} \]

Further, define the condition C:
\[ C: K_{XY} = \arg \min_G \log \left| \begin{bmatrix} I & 0 \\ 0 & K_{XY}K_Y^{-1} \end{bmatrix} K \begin{bmatrix} I & 0 \\ 0 & K_{XY}K_Y^{-1} \end{bmatrix} \right| \tag{11} \]
\[ -\log |\Gamma| + \text{Tr} \left\{ \left( \begin{bmatrix} I & 0 \\ 0 & K_{XY}K_Y^{-1} \end{bmatrix} K \begin{bmatrix} I & 0 \\ 0 & K_{XY}K_Y^{-1} \end{bmatrix} \right) \Gamma \right\} \]
where the minimum is over all $m$-by-$q$ matrices $G$ such that the matrix
\[ \Gamma \triangleq \begin{bmatrix} K_X & G \bar{K}_{XY} \bar{K}_Y^{-1}K_{XY} \bar{K}_{XY}^{-1}K_Y \end{bmatrix}. \tag{12} \]
is positive semi-definite, i.e.,
\[ \Gamma \succeq 0. \tag{13} \]
The following theorem provides an explicit analytic expression of the exponent-rate function of the vector Gaussian hypothesis testing problem of Figure 1 when condition C in (11) is fulfilled.

**Theorem 1.** If C is satisfied, then
\[ E(R) = \frac{m}{2} + \frac{q}{2} + \frac{1}{2} \log |\bar{K}_Y| + \frac{1}{2} \text{Tr} (\bar{K}_X^{-1}K_Y) \]
\[ + \frac{1}{2} \log |\bar{K}_X| - \log |K_X - K_{XY}\bar{K}_Y^{-1}K_{XY}^\dagger| \]
\[ + \frac{1}{2} \text{Tr} (\bar{K}_X^{-1}K_XK_YK_Y^{-1}K_{XY}^\dagger) \]
\[ + \frac{1}{2} \text{Tr} (\bar{K}_X^{-1}K_{XY}K_YK_Y^{-1}K_{XY}^\dagger) \]
\[ + \max \min \left\{ R + \frac{1}{2} \log |I - \Omega K_X|, \right. \]
\[ \left. \frac{1}{2} \log |I + \Omega K_X| \right\} \tag{14} \]
where the maximization in the last term is over all matrices $0 \leq \Omega \leq K_{XY}^\dagger$ and where $K_{XY}^\dagger$ designates the Moore-Penrose pseudo inverse of $K_{XY}$.

**Proof:** See Section IV.

**Remark 1.** The theorem recovers the result of [4, Theorem 7] in the special case of testing against independent and $m = q = 1$. In this case, when the distribution $P_{XY}$ under the null hypothesis $H = 0$ describes the channel $Y = X + N$ with $X$ and $N$ independent Gaussian both with zero mean and respective variances $\sigma_X^2$ and $\sigma_N^2$, and the joint law $P_{XY}$ under $H = 1$ describes a pair of independent Gaussians of variances $\sigma_X^2 + \sigma_N^2$, and $\sigma_N^2$, then:
\[ E(R) = \frac{1}{2} \log \left( \frac{\sigma_X^2 + \sigma_N^2}{\sigma_N^2 + 2 - 2R\sigma_X^2} \right). \tag{15} \]

### IV. Proof of Theorem

We first derive an auxiliary result. Consider a slight generalization of the discrete memoryless single-sensor single-detector hypothesis testing against conditional independence problem where the marginals are not identical under the two hypotheses. Specifically, consider the problem of Figure 2 where under
\[ H = 0: \quad \{(X_i, U_i, V_i)\}_{i=1}^n \text{ i.i.d. } \sim P_{XUV} \tag{16a} \]
\[ P_{XU} = \arg\min_{\bar{P}_{XU}: \bar{P}_{X} = \bar{P}_{X}, \bar{P}_{U} = P_{U}} D(\bar{P}_{XU} || \bar{P}_{XU}), \] (17)

the rate exponent function is given by

\[ E(R) = D(P_{XU} || \bar{P}_{XU}) + E_{P_{V}}[D(P_{V} || \bar{P}_{V})] + \max I(S; V | U) \] (18)

where in (18) the maximization is over all conditionals p.m.f.s \( P_{S|X} \) for which \( I(S; X | U) \leq R \).

Proof of Lemma 1: By (1) Theorem 4:

\[ E(R) = \lim_{n \to \infty} E_{n}(R), \] (19)

where

\[ E_{n}(R) \triangleq \max_{\phi_{n}} \frac{1}{n} D(\phi_{n}(X^{n})U^{n}V^{n} || \bar{P}_{\phi_{n}(X^{n})U^{n}V^{n}}). \] (20)

Next, notice that by the chain rule for KL divergence, the data processing inequality, and some simple manipulations, we have

\[ D(\phi_{n}(X^{n})U^{n}V^{n} || \bar{P}_{\phi_{n}(X^{n})U^{n}V^{n}}) \]
\[ = D(\phi_{n}(X^{n})U^{n} || \bar{P}_{\phi_{n}(X^{n})U^{n}}) + E_{P_{\phi_{n}(X^{n})U^{n}}}[D(P_{V} || \bar{P}_{V})] + I(V^{n}; \phi_{n}(X^{n})|U^{n}) + nE_{P_{U}}[D(P_{V} || \bar{P}_{V})] \] (21)
\[ \leq nD(P_{XU} || \bar{P}_{XU}) + I(V^{n}; \phi_{n}(X^{n})|U^{n}) + nE_{P_{U}}[D(P_{V} || \bar{P}_{V})], \] (22)

where (a) holds by the data-processing inequality for KL divergence and because \( X^{n} \) and \( U^{n} \) are i.i.d.

We can thus bound \( E(R) \) as:

\[ E(R) \leq D(P_{XU} || \bar{P}_{XU}) + E_{P_{V}}[D(P_{V} || \bar{P}_{V})] + \lim_{n \to \infty} \max_{\phi_{n}} \frac{1}{n} I(V^{n}; \phi_{n}(X^{n})|U^{n}). \] (24)

Next we use that by (1) Theorem 4 and (4) Theorem 3 both sides of

\[ \lim_{n \to \infty} \max_{\phi_{n}} \frac{1}{n} \cdot (I(\phi_{n}(X^{n}); V^{n}|U^{n})) = \max_{P_{S|X}} I(S; V | U) \] (25)

characterize the optimal type-II error exponent of a hypothesis testing against conditional independence problem at rate \( R \), and thus coincide.

Combining (19) and (25) we obtain:

\[ E(R) \leq D(P_{XU} || \bar{P}_{XU}) + E_{P_{V}}[D(P_{V} || \bar{P}_{V})] + \max I(S; V | U). \] (26)

The reverse inequality follows from the achievable type-II error exponent of Shimokawa-Han-Amari (SHA) (see Section IV) for an analysis) which states that for every choice of the conditional distribution \( P_{S|X} \) satisfying \( R \geq I(S; X | U) \) the following lower bound holds:

\[ E(R) \geq \min \left\{ \min_{P_{S|X}} D(\bar{P}_{S|X} || P_{S|X} \bar{P}_{XU} \bar{P}_{V}), \min_{P_{S|X}} D(\bar{P}_{S|X} || P_{S|X} \bar{P}_{XU} \bar{P}_{V}) \right\} \]
\[ + R - I(S; X | U) \] (27)

where the mutual information \( I(S; X | U) \) is calculated according to \( P_{S|X} P_{U} P_{V} \). In what follows we show that the SHA result implies that the error exponent on the RHS of (18) is achievable. As in (4), we restrict to distributions \( P_{S|X} \) satisfying

\[ R \geq I(S; X | U) \] (28)

and drop the condition \( H(S; U, V) \leq H_{S|X} P_{U} P_{V} \). These changes can lead to a smaller exponent than in [4], and thus the resulting exponent is still achievable.

By the Markov chain \( S \rightarrow X \rightarrow (U, V) \), Condition (28) implies

\[ R - I(S; X | U) \geq I(S; V | U). \] (29)

Moreover, by the chain rule and the nonnegativity and convexity of KL divergence, for any \( \bar{P}_{S|X} \bar{P}_{XU} \bar{P}_{V} \):

\[ D(\bar{P}_{S|X} || P_{S|X} \bar{P}_{XU} \bar{P}_{V}) \geq D(\bar{P}_{XU} || P_{XU} \bar{P}_{V}) \]
\[ \geq D(P_{XU} || \bar{P}_{XU}) + E_{P_{V}}[D(\bar{P}_{V} || \bar{P}_{V})] \] (30)
\[ \geq D(\bar{P}_{XU} || \bar{P}_{XU}) + E_{P_{U}}[D(\bar{P}_{V} || \bar{P}_{V})]. \] (31)

By (29) and (30) and since the second minimization in (27) is over distributions \( P_{S|X} \) satisfying \( P_{U} = P_{V} \), we conclude that under conditions (17) and (28) the second term in (27) is lower bounded by

\[ \theta \triangleq D(\bar{P}_{XU} || \bar{P}_{XU}) + E_{P_{V}}[D(\bar{P}_{V} || \bar{P}_{V})] + I(S; V | U). \] (33)
We now lower bound the first term in (27). By the chain rule and the nonnegativity and convexity of KL divergence, for any \( \hat{P}_{SUV} \) where \( \hat{P}_{SX} = P_{S|X} P_X \):
\[
D(\hat{P}_{SUV} \| P_{S|X} \hat{P}_X \hat{P}_V) \\
\quad = D(\hat{P}_{SX} \| \hat{P}_X) + \mathbb{E}_{P_{S|X}}[D(\hat{P}_{SV|XU} \| P_{S|X} \hat{P}_V)] \\
\quad \geq D(\hat{P}_{SX} \| \hat{P}_X) + \mathbb{E}_{P_X}[D(\hat{P}_{SV|U} \| P_{S|X} \hat{P}_V)]. \tag{34}
\]

Where in the last inequality we used that \( \sum_x \hat{P}_X(x) P_{S|X}(s|x) = P_S(s) \) because \( \hat{P}_X = P_X \). We now notice that the first minimization in (27) is only over distributions \( \hat{P}_{SUV} \) satisfying \( \hat{P}_{S|X} = P_{S|X} \) and therefore:
\[
\mathbb{E}_{P_X}[D(\hat{P}_{SV|U} \| P_{S} \hat{P}_V)] \\
\quad = \mathbb{E}_{P_X}[D(\hat{P}_{SV|U} \| P_{S} \hat{P}_V)] \\
\quad = I(S; V|U) + \mathbb{E}_{P_X}[D(\hat{P}_{SV|U} \| P_{S} \hat{P}_V)]. \tag{35}
\]

Combining (34) and (35) we conclude that under Condition 17 also the first term in the minimization in (27) is lower bounded by \( \theta \). This establishes the achievability of the right-hand side of (38).

We turn to the proof of Theorem 1 Define
\[
U = \mathbb{E}_P[X|Y] \tag{37}
\]
\[
V = Y \tag{38}
\]
and notice that under \( H = 1 \) they satisfy the Markov chain
\[
X \rightarrow U \rightarrow V. \tag{39}
\]

In what remains, we assume that instead of \( Y^n \) the decoder observes the pair of sequences \((U^n, V^n)\) which is i.i.d according to the joint distribution of \((U, V)\). This new system is depicted in Figure 2. Since there is a bijection between \( Y^n \) and \((U^n, V^n)\), the error exponent of the new system coincides with the error exponent of the original system. Moreover, by the Markov chain (39) the new system is a generalized testing against conditional independence problem as described in (16).

We next argue that under condition C in Theorem 1 and because \( \hat{P}_{X^n|U^n} \) and \( \hat{P}_{X^n|U^n} \) are multivariate Gaussian distributions, the new system also satisfies Condition 17 in Lemma 1. The optimal exponent \( E(R) \) will then follow immediately from this Lemma 1. To show that for multivariate Gaussian distributions \( \hat{P}_{X^n|U^n} \) and \( \hat{P}_{X^n|U^n} \) condition C in (11) implies (17), we first show that under this Gaussian assumption the minimizer of
\[
\arg \min_{\hat{P}_{S|U; \hat{P}_X = P_X; \hat{P}_U = P_U}} D(\hat{P}_{SU} \| \hat{P}_{SU}) \tag{40}
\]
is a multivariate Gaussian distribution. To see this fix any distribution \( \hat{P}_{SU} \) with \( \hat{P}_X = P_X \) and \( \hat{P}_U = P_U \) and let \( \hat{P}_{SU} \) be a multivariate Gaussian distribution with covariance matrix as \( \hat{P}_{SU} \). Then:
\[
D(\hat{P}_{SU} \| \hat{P}_{SU}) = -h(\hat{P}_{SU}) - \mathbb{E}_{\hat{P}}[\log \hat{P}_{SU}] \\
\quad \geq -h(\hat{P}_{SU}) - \mathbb{E}_{\hat{P}}[\log \hat{P}_{SU}] \\
\quad = -h(\hat{P}_{SU}^G - \mathbb{E}_{\hat{P}}[\log \hat{P}_{SU}], \tag{41}
\]
where the inequality holds because a Gaussian distribution maximizes differential entropy under a fixed covariance matrix constraint and where the last equality holds because
\[\mathbb{E}[\log \hat{P}_{SU}] \]
only depends on the covariance matrix of \((U, X)\) which is the same under \( P \) and \( P^G \). By straightforward algebra, it can then be shown that if condition C in (11) holds, then \( \hat{P}_{SU} \) is the multivariate Gaussian distribution that minimizes (17).

We conclude that the optimal exponent \( E(R) \) is given by (18) in Lemma 1. We evaluate (18) for our problem. For simplicity, we rewrite
\[
D(\hat{P}_{SU} \| \hat{P}_{SU}) + \mathbb{E}_{P_{U|V}}[D(\hat{P}_{V|U} \| \hat{P}_{V|U})] \\
\quad = D(\hat{P}_{UV} \| \hat{P}_{UV}) + \mathbb{E}_{P_{X|U}}[D(\hat{P}_{X|U} \| \hat{P}_{X|U})]. \tag{42}
\]
and proceed to compute
\[
D(\hat{P}_{UV} \| \hat{P}_{UV}) = \frac{m}{2} + \frac{1}{2} \log |\hat{K}_Y| + \frac{1}{2} \text{Tr} (\hat{K}_Y^{-1} \hat{K}_Y) \tag{43}
\]
and
\[
D(\hat{P}_{X|U} \| \hat{P}_{X|U}) = \frac{m}{2} + \frac{1}{2} \log |\hat{K}_X| - \frac{1}{2} \log |\hat{K}_X - \hat{K}_{XY} \hat{K}_Y^{-1} \hat{K}_{XY}^\dagger| + \frac{1}{2} \text{Tr} (\hat{K}_X^{-1} \hat{K}_Y^{-1} \hat{K}_{XY} \hat{K}_Y^{-1} \hat{K}_{XY}^\dagger)^+ \\
\quad \times \hat{K}_{XY} \hat{K}_Y^{-1} \hat{K}_X \hat{K}_Y^{-1} \hat{K}_{XY} \hat{K}_Y^{-1} \hat{K}_{XY} \tag{44}
\]

It remains to find \( \max I(S; Y|U) \) where the maximum is over all test channels \( P_{S|X} \) satisfying \( I(S; X|U) \leq R \). Let \( \hat{U} = U + \epsilon Z \) where \( Z \sim (0, I) \). Applying the result of [3] Theorem 2) on the triple \((X, Y, \hat{U})\), which is Gaussian, and then taking the limit \( \epsilon \rightarrow 0 \) we get:
\[
\max I(S; Y|U) = \max \min \left\{ R + \frac{1}{2} \log |I - \Omega \hat{K}_X| : \right. \\
\left. I(S; X|U) \leq R \right\} \\
\quad = \frac{1}{2} \log |I + \Omega \hat{K}_X| \\
\quad \times \left( \hat{K}_X^{-1} - \hat{K}_Y^{-1} \hat{K}_{XY} \hat{K}_Y^{-1} \hat{K}_{XY} \hat{K}_X \hat{K}_Y^{-1} \hat{K}_{XY} \right)^{1/2}, \tag{45}
\]
where the maximization in the last term is over all matrices \( 0 \leq \Omega \leq \hat{K}_X^{-1} \) and where \( \hat{K}_X \) designates the Moore-Penrose pseudo inverse of \( \hat{K}_X \).

Summing (43) - (45) we obtain the desired result in (44), which completes the proof of Theorem 1.

V. DISCUSSION

In what follows, we show that constraint C as given by (11) is fulfilled for a large class of sources even when \( m = 1 \) and \( q = 2 \). Let \( X \) be a scalar source that is observed at the sensor and \( Y = (Y_1, Y_2) \) a 2-dimensional source that is observed at the detector. For convenience, let
\[
\hat{K} = \begin{bmatrix}
\sigma_X^2 & \sigma_{XY_1} & \sigma_{XY_2} \\
\sigma_{XY_1} & \sigma_{Y_1}^2 & \sigma_{Y_1Y_2} \\
\sigma_{XY_2} & \sigma_{Y_1Y_2} & \sigma_{Y_2}^2
\end{bmatrix}
\]
and
\[
\hat{K} = \begin{bmatrix}
\sigma_X^2 & \sigma_{XY_1} & \sigma_{XY_2} \\
\sigma_{XY_1} & \sigma_{Y_1}^2 & \sigma_{Y_1Y_2} \\
\sigma_{XY_2} & \sigma_{Y_1Y_2} & \sigma_{Y_2}^2
\end{bmatrix}. \tag{46}
\]
Also, let
\[ a = (\bar{\sigma}_{XY}^2\bar{\sigma}_{YZ}^2 - \bar{\sigma}_{XY}\bar{\sigma}_{YZ}^2) \quad \text{and} \quad b = (\bar{\sigma}_{XY}^2\bar{\sigma}_{YZ}^2 - \bar{\sigma}_{XY}\bar{\sigma}_{YZ}^2). \]
For example, the constraint \( C \) as given by (11) reduces to
\[
\begin{align*}
\text{i)} \quad & \sigma_X^2 = \sigma_Y^2, \\
\text{ii)} \quad & a(\sigma_{XY} - \sigma_{XY}) + b(\sigma_{XZ} - \sigma_{XZ}) = 0, \\
\text{iii)} \quad & a^2(\sigma_Y^2 - \bar{\sigma}_Y^2) + 2ab(\sigma_{YZ} - \sigma_{YZ}) + b^2(\sigma_Z^2 - \bar{\sigma}_Z^2) = 0.
\end{align*}
\]

For example, if all components have unit variance under both \( P \) and \( P_r \), i.e., \( \sigma_X^2 = \sigma_Y^2 = \sigma_Z^2 = 1 \) and \( \sigma_{XY}^2 = \sigma_{YZ}^2 = \sigma_{XZ}^2 = \bar{\sigma}_Y^2 = 1 \) then all definite positive matrices \( \mathbf{K} \) and \( \mathbf{K} \) of the form
\[
\mathbf{K} = \begin{bmatrix}
1 & a_{12} & h(\bar{a}_{12}, \bar{a}_{13}, \bar{a}_{23}, a_{12}) \\
\bar{a}_{12} & 1 & \bar{a}_{23} \\
\bar{a}_{13} & \bar{a}_{23} & 1
\end{bmatrix},
\]
and
\[
\mathbf{\bar{K}} = \begin{bmatrix}
1 & a_{12} & \bar{a}_{13} \\
\bar{a}_{12} & 1 & \bar{a}_{23} \\
\bar{a}_{13} & \bar{a}_{23} & 1
\end{bmatrix}
\]
for some arbitrary parameters \( a_{12}, \bar{a}_{12}, \bar{a}_{13}, \bar{a}_{23} \), satisfy the constraint \( (47) \). Here
\[
h(x, y_1, y_2, t) = y_1 - (t - y_2) = \frac{y_1y_2 - x}{x^2y_2 - y_1}.
\]

\section*{Example 1.}
Let
\[
\mathbf{K} = \begin{bmatrix}
1 & 0.4 & \alpha \\
0.4 & 1 & 0.1 \\
\alpha & 0.1 & 1
\end{bmatrix}
\quad \text{and} \quad
\mathbf{\bar{K}} = \begin{bmatrix}
1 & 0.1 & -0.8 \\
0.1 & 1 & 0.1 \\
-0.8 & 0.1 & 1
\end{bmatrix},
\]
with \( \alpha \approx -0.73333 \). It is easy to see that \( (47) \) is fulfilled. Figure 3 shows the evolution of the optimal exponent \( E \) as a function of the communication rate \( R \) as given by Theorem 7 for this example. Notice that Han’s exponent \( [2 \text{ Theorem } 2] \) is strictly suboptimal for this example.

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\begin{figure}[h]
\centering
\includegraphics[width=0.5\textwidth]{rate_exponent_region.png}
\caption{Rate-exponent region for Example 1}
\end{figure}

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