I know what you mean: semantic issues in Internet-scale publish/subscribe systems

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Abstract. In recent years, the amount of information on the Internet has increased exponentially developing great interest in selective information dissemination systems. The publish/subscribe paradigm is particularly suited for designing systems for routing information and requests according to their content throughout wide-area network of brokers. Current publish/subscribe systems use limited syntax-based content routing but since publishers and subscribers are anonymous and decoupled in time, space and location, often over wide-area network boundary, they do not necessarily speak the same language. Consequently, adding semantics to current publish/subscribe systems is important. In this paper we identify and examine the issues in developing semantic-based content routing for publish/subscribe broker networks.

1 Introduction

The increase in the amount of data on the Internet has led to the development of a new generation of applications based on selective information dissemination where data is distributed only to interested clients. Such applications require a new middleware architecture that can efficiently match user interests with available information. Middleware that can satisfy this requirement include event-based architectures such as publish/subscribe systems.

In publish/subscribe systems (hereafter referred to as pub/sub systems), clients are autonomous components that exchange information by publishing events and by subscribing to the classes of events they are interested in. In these systems, publishers produce information, while subscribers consume it. A component usually generates a message when it wants the external world to know that a certain event has occurred. All components that have previously expressed their interest in receiving such events will be notified about it. The central component of this architecture is the event dispatcher (also known as event broker). This component records all subscriptions in the system. When a certain event is published, the event dispatcher matches it against all subscriptions in the system. When the incoming event verifies a subscription, the event dispatcher sends a notification to the corresponding subscriber.

1 We use the terms event and publication interchangeable.
The earliest pub/sub systems were subject-based. In these systems, each message (event) belongs to a certain topic. Thus, subscribers express their interest in a particular subject and they receive all the events published within that particular subject. The most significant restriction of these systems is the limited selectivity of subscriptions. The latest systems are called content-based systems. In these systems, the subscriptions can contain complex queries on event content.

Pub/sub systems try to solve the problem of selective information dissemination. Recently, there has been a lot of research on solving the problem of efficiently matching events against subscriptions. The proposed solutions are either centralized, where a single broker stores all subscriptions and event matching is done locally [1–3] or distributed, where many brokers need to collaborate to match events with subscriptions because not all subscriptions are available to every broker [4, 5]. The latter approach is also referred to as content-based routing because brokers form a network where events are routed to interested subscribers based on their content.

The existing solutions are limited because the matching (routing) is based on the syntax and not on the semantics of the information exchanged. For example, someone interested in buying a car with a “value of up to 10,000 will not receive notifications about “vehicles, “automobiles or even “cars with “price of 8,999 because the system has no understanding of the “price-“value relationship, nor of the “car-“automobile-“vehicle relationship.

In this paper we examine the issues in extending distributed pub/sub systems to offer semantic capabilities. This is an important aspect to be studied as components in a pub/sub systems are decoupled and do not necessary speak the same language.

2 Local Matching and Content-based Routing

Due to space limitation, we will not provide an extensive background about pub/sub systems and content-based routing. Instead, we briefly present the most important concepts that help the reader understand the ideas conceived in this paper.

The key point in pub/sub systems is that the information sent into the system by the publisher does not contain the addresses of the receivers. The information is forwarded to interested clients based on the content of the message and clients subscriptions. In a centralized approach, there is only one broker that stores all subscriptions. Upon receiving an event, the broker uses a matching algorithm to match the event against the subscriptions in order to decide which subscribers want to receive notifications about the event.

Usually, publications are expressed as lists of attribute-value pairs. The formal representation of a publication is given by the following expression: \((a_1, \text{val}_1), (a_2, \text{val}_2), \ldots, (a_n, \text{val}_n)\). Subscriptions are expressed as conjunctions of simple predicates. In a formal description, a simple predicate is represented as \((\text{attribute name relational operator value})\). A predicate \((a \text{rel} \text{op} \text{val})\) is matched by an attribute-value pair \((a, \text{val})\) if and only if the attribute names are identi-
cal \( (a = a) \) and the \( (a \text{ relOp} val) \) boolean relation is true. A subscription \( s \) is matched by a publication \( p \) if and only if all its predicates are matched by some pair in \( p \). In this case we say that the subscription is matched at syntactic level.

The distributed approach involves a network of brokers that collaborate in order to route the information in the system based on its content. In this case, practically, each broker is aware of its neighbours interests. Upon receiving an event, the broker matches it against its neighbours subscriptions and sends the event only to the interested neighbours. Usually, the routing scheme presents two distinct aspects: subscription forwarding and event forwarding. Subscription forwarding is used to propagate clients interests in the system, while event forwarding algorithms decide how to disseminate the events to the interested clients. Two main optimizations were introduced in the literature in order to increase the performance of these forwarding algorithms: subscription covering and advertisements.

**Subscription covering**

Given two subscriptions \( s_1 \) and \( s_2 \), \( s_1 \) covers \( s_2 \) if and only if all the events that match \( s_2 \) also match \( s_1 \). In other words, if we denote with \( E_1 \) and \( E_2 \) the set of events that match subscription \( s_1 \) and \( s_2 \), respectively, then \( E_2 \subseteq E_1 \).

If we look at the predicate level, the covering relation can be expressed as follows: Given two subscriptions \( s_1 = p_{11}, p_{21}, \ldots, p_{n1} \) and \( s_2 = p_{12}, p_{22}, \ldots, p_{m2} \), \( s_1 \) covers \( s_2 \) if and only if \( \forall p_{k1} \in s_1, \exists p_{j2} \in s_2 \) such that if \( p_{j2} \) is matched by some attribute-value pair \((a, val)\), then \( p_{k1} \) is also matched by the same \((a, val)\) attribute-value pair. In other words, \( s_2 \) has potentially more predicates and they are more restrictive than those in \( s_1 \). Table 1 presents some examples of subscriptions and the corresponding covering relations.

When a broker \( B \) receives a subscription \( s \), it will send it to its neighbours if and only if it has not previously sent them another subscription \( s' \), that covers \( s \). Broker \( B \) is ensured to receive all events that match \( s \), since it receives all events that match \( s \) and the events that match \( s \) are included in the set of the events that match \( s \).

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| Subscription \( s_1 \) | Subscription \( s_2 \) | Covering Relation |
|------------------------|------------------------|------------------|
| \((\text{product} = \text{"computer, brand} = \text{"IBM, price} \leq 1600)\) | \((\text{product} = \text{"computer, brand} = \text{"IBM, price} \leq 1500)\) | \( s_1 \text{ covers } s_2 \) |
| \((\text{product} = \text{"computer, brand} = \text{"IBM, price} \leq 1600)\) | \((\text{product} = \text{"computer, price} \leq 1600)\) | \( s_2 \text{ covers } s_1 \) |
| \((\text{product} = \text{"computer, brand} = \text{"IBM, price} \leq 1600)\) | \((\text{product} = \text{"computer, brand} = \text{"Dell, price} \leq 1500)\) | \( s_1 \text{ does not cover } s_2, s_2 \text{ does not cover } s_1 \) |

Table 1. Examples of subscriptions and covering relations

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\( p_{k1} \) and \( p_{j2} \) refer to the same attribute
Table 2. Examples of subscriptions, advertisements and intersection relations

| Subscription s | Advertisement a | Intersection Relation |
|----------------|----------------|-----------------------|
| (product = “computer, brand = "IBM, price ≤ 1600) | (product = “computer, brand = "IBM, price ≤ 1500) | a intersects s |
| (product = “computer, price ≤ 1600) | (product = “computer, brand = "IBM, price ≤ 1600) | a intersects s |
| (product = “computer, brand = "IBM, price ≤ 1600) | (product = “computer, brand = "Dell, price ≤ 1500) | a does not intersect s |

Advertisements

Advertisements are used by publishers to announce the set of publications they are going to publish. Advertisements look exactly like subscriptions, but have a different role in the system: they are used to build the routing path from the publishers to the interested subscribers.

An advertisement \(a\) determines an event \(e\) if and only if all attribute-value pairs match some predicates in the advertisement. Formally, an advertisement \(a = p_1^1, p_2^1, \ldots, p_n^1\) determines an event \(e\), if and only if \(\forall (a, v) \in e, \exists p_k \in a\) such that \((a, v)\) matches \(p_k\).

An advertisement \(a\) intersects a subscription \(s\) if and only if the intersection of the set of the events determined by the advertisement \(a\) and the set of the events that match \(s\) is a non-empty set. Formally, at predicate level, an advertisement \(a = a_1, a_2, \ldots, a_n\) intersects a subscription \(s = s_1, s_2, \ldots, s_n\) if and only if \(\forall s_k \in s, \exists a_j \in a\) and some attribute-value pair \((\text{attr}, \text{val})\) such that \((\text{attr}, \text{val})\) matches both \(s_k\) and \(a_j\). Table 2 presents some examples of subscriptions and advertisements and the corresponding intersection relations.

When using advertisements, upon receiving a subscription, each broker forwards it only to the neighbours that previously sent advertisements that intersect with the subscription. Thus, the subscriptions are forwarded only to the brokers that have potentially interesting publishers.

3 Towards Semantic-based Routing

In order to add a semantic dimension to distributed pub/sub systems, we have to understand how to adapt or map the core concepts and functionalities of existing solutions for content-based routing to the new context that involves semantic knowledge.

\(^3\) However, there is an important distinction between the predicates in an advertisement and those in a subscription: the predicate in a subscription are considered to be in a conjunctive construction, while those in an advertisement behave as in a disjunctive one.

\(^4\) \(s_k\) and \(a_j\) refer to the same attribute \(\text{attr}\).
In this section we first introduce some extensions to the existing matching algorithms in order to make them semantic-aware and then we discuss the implications of using such a solution for semantic-based routing.

### 3.1 Semantic Matching

In this section we summarize our approach to make the existing centralized matching algorithms semantic-aware [6]. Our goal is to minimize the changes to the existing matching algorithms so that we can take advantage of their already efficient techniques and to make the processing of semantic information fast. We describe three approaches, each adding more extensive semantic capability to the matching algorithms. Each of the approaches can be used independently and for some applications that may be desirable. It is also possible to use all three approaches together.

The first approach allows a matching algorithm to match events and subscriptions that use semantically equivalent attributes or values—synonyms. The second approach uses additional knowledge about the relationships (beyond synonyms) between attributes and values to allow additional matches. More precisely, it uses a concept hierarchy that provides two kinds of relations: specialization and generalization. The third approach uses mapping functions which allow definitions of arbitrary relationships between the schema and the attribute values of the event.

The synonym step involves translating all strings with different names but with the same meaning to a “root term. For example, “car” and “automobile” are synonyms for “vehicle” which then becomes the root term for the three words. This translation is performed for both subscriptions and events and at both attribute and value level. This allows syntactically different events and subscriptions to match. This translation is simple and straightforward. The semantic capability it adds to the system, although important, may not be sufficient in some situations, because this approach operates only at attribute and value level independently and does not consider the semantic relation between attributes and values. Moreover, this approach is limited to synonym relations only.

Taxonomies represent a way of organizing ontological knowledge using specialization and generalization relationships between different concepts. Intuitively, all the terms contained in such a taxonomy can be represented in a hierarchical structure, where more general terms are higher up in the hierarchy and are linked to more specialized terms situated lower in the hierarchy. This structure is called a “concept hierarchy. Usually, a concept hierarchy contains all terms within a specific domain, which includes both attributes and values.

Considering the observation that the subscriber should receive only information that it has precisely requested, we come up with the following two rules for matching that uses concept hierarchy: (1) the events that contain more specialized concepts have to match the subscriptions that contain more generalized terms of the same kind and (2) the events that contain more generalized terms than those used in the subscriptions do not match the subscriptions.
In order to better understand these rules, we look at the following examples. Suppose that we have in the system a subscription:

\[ S : (book = Stone.Age) AND (subject = reptiles). \]

When the event:

\[ E : (encyclopedia, Stone.Age), (subject, crocodiles) \]

is entering the system, it should match the subscription \( S \), as the subscriber asked for more general information that the event provides (in other words, an encyclopedia is a special kind of book and crocodiles represent a special kind of reptiles). On the other hand, considering the subscription:

\[ S : (encyclopedia = Stone.Age) AND (subject = reptiles) \]

and the incoming event

\[ E : (book, Stone.Age), (subject, crocodiles) \]

the event \( E \) should not match the subscription \( S \), as the book contained in the event may be a dictionary or a fiction book (as well as an encyclopedia). Note that, although the subscription \( S \) contains in its second predicate a value more specialized than that in the event, the first predicate of the subscription is not matched by the event, and therefore, the event does not match the subscription. The last rule prevents an eventual spamming of the subscribers with useless information.

Mapping functions can specify relationships which otherwise cannot be specified using a concept hierarchy or a synonym relationship. For example, they can be used to create a mapping between different ontologies. A mapping function is a many-to-many function that correlates one or more attribute-value pairs to one or more semantically related attribute-value pairs. It is possible to have many mapping functions for each attribute. We assume that mapping functions are specified by domain experts. In the future, we are going to investigate using a fully-fledged inference engine as a more compact representation of mapping functions and the performance trade off this entails.

We illustrate the concept of mapping functions with an example. Let us say that there is a university professor X, who is interested in advising new PhD graduate students. In particular, he is only interested in students who have had 5 or more years of previous professional experience. Subsequently, he subscribes to the following:

\[ S : (university = Y) AND (degree = PhD) AND (professional experience > 4) \]

Specifically, the professor X is looking for students applying to university \( Y \) in the PhD stream with 5 or more years of experience. For each new student applying to the university, a new event, which contains (among others) the information about previous work experience, is published into our system. Thus, an event for a student who had some work experience would look like

\[ E : (school, Y)(degree, PhD)(work experience, true)(graduation date, 1990). \]

In addition, the system has access to the following mapping function:

\[ f_1 : (work experience, graduation date) \rightarrow professional experience. \]

You can think of function \( f_1 \) implemented as a simple difference between today's date and the date of students graduation and returning that difference as the value of professional experience. For the purposes of the example, \( f_1 \) assumes
that the student has been working since graduation. Finally, the result of $f_1$ is appended to event $E$ and the matching algorithm matches $E$ to professor X’s subscription $S$.

In addition, we can think about events and subscriptions as points in a multi-dimensional space [7] where the distance between points determines a match between an event and a subscription. This way it is possible that an event matches a subscription even if some attribute/value pair of the event is more general than the corresponding predicate in the subscription as long as the distance between the event and the subscription, as determined by all their constituent attribute-value pairs and predicates respectively, is within the defined matching range.

To summarize, the synonym stage translates the events and the subscriptions to a normalized form using the root terms, while the hierarchy and the mapping stages add new attribute-value pairs to the events. The new events are matched using existing matching algorithms against the subscriptions in the system. In conclusion, we say that $e$ semantically matches $s^5$ if and only if the hierarchy and the mapping stages can produce an event $e = e \cup E^6$ that matches $s$ at syntactic level.

3.2 Semantic-based Routing

At first glance, it is apparent that existing algorithms for subscription and event forwarding can be used with a semantic-aware matching algorithm in order to achieve semantic-based routing. However, this approach is not straight forward. In this section we discuss some open issues that arise from using a semantic-aware matching algorithm in content-based routing.

**Subscription covering**

Although it is defined at syntax level, the covering relation as presented in Section 2 can be used directly with the semantic matching approach discussed above without any loss of notifications. In other words, if $s_1$ covers $s_2$ and a certain broker $B$ will forward only subscription $s_1$ to its neighbours, it will still receive both events that semantically match $s_1$ and $s_2$. This happens because the relation between the set of events $E_1$ and $E_2$ that semantically match $s_1$ and $s_2$ respectively, is preserved, i.e. $E_2 \subseteq E_1$. Truly, if $e$ semantically matches $s_2$, then the hierarchy and the mapping stages can produce an event $e$ that matches $s_2$ at syntactic level. If $e$ matches $s_2$ at syntactic level, then, according to the definition of covering relation, $e$ matches $s_1$ at syntactic level. Since $e$ is produced by adding semantic knowledge to $e$, this means that $e$ semantically matches $s_1$, i.e. $E_2 \subseteq E_1$. Thus, broker $B$ is ensured to receive all events that semantically match $s_2$, since it receives all events that semantically match $s_1$ and the events that semantically match $s_2$ are included in the set of the events that semantically match $s_1$.

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5 $e$ and $s$ are in their normalized form

6 $E$ represents the set of attribute-value pairs that are added by the hierarchy and the mapping stages. Note that $E$ can be an empty set.
Although the syntactic covering relation can be used without loss of notifications, some redundant subscriptions may be forwarded into the network. This happens because the set of events \( E_1 \) and \( E_2 \) that semantically match \( s_1 \) and \( s_2 \) can be in the following relation \( E_2 \subseteq E_1 \) without necessarily \( s_1 \) covering \( s_2 \) at syntax level. In other words, although \( s_1 \) does not cover \( s_2 \) at syntactic level, it may cover it semantically speaking. For example, consider the following subscriptions: \( s_1 = (\text{product} = \text{"printed material"}, \text{AND(topic} = \text{"semantic web"})) \) and \( s_2 = (\text{product} = \text{"book"}, \text{AND(topic} = \text{"semantic web"})) \). In this case, all events that semantically match \( s_2 \) will also match \( s_1 \) as a book is a form of printed material; thus \( E_2 \subseteq E_1 \), but \( s_1 \) does not cover \( s_2 \) (at syntax level). Therefore, the covering relation needs to be extended to encapsulate semantic knowledge. One simple way of transforming the covering relation to be semantic-aware is to use the hierarchy approach. In this case, subscription \( s_1 \) will cover \( s_2 \) as the printed material term is a more general term than book.

Advertisements

While the covering relation can be directly used with the semantic matching algorithms, this is not the case for advertisements. As explained earlier in this paper, advertisements are used to establish the routing path from the publishers to the interested subscribers. How the events are routed in the system depends on the intersection relation between advertisements and subscription. Consider the following example: advertisement \( a = (\text{product} = \text{"printed material"}, (\text{price} \geq 10)) \) and subscription \( s = (\text{product} = \text{"book"}, (\text{price} \leq 20)) \). Advertisement \( a \) does not intersect \( s \) at syntactic level because there is not any predicate \( p \) in \( a \) and not any attribute-value pair \((\text{attr}, \text{val})\) such that \((\text{attr}, \text{val})\) matches both \( p \) and the following predicate \((\text{product} = \text{"book"})\) of subscription \( s \). (cf. Section 2. Thus, the subscription will not be forwarded towards the publisher that emitted the advertisement. All publications that will be produced by this publisher will not be forwarded to the subscriber, although some of them may matched its subscriptions.

Distributed semantic knowledge

The discussion above about subscription covering and advertisements considered that each broker contains the same semantic knowledge (i.e. same synonyms, hierarchies and mapping functions). However, the replication of the same semantic knowledge to all brokers in the system may not be feasible and it may be detrimental to scalability.

We envision a system where semantic knowledge is distributed between brokers\(^7\) in the same way that Internet distributes link status information using routing protocols. A semantic knowledge database is equivalent to routing tables in terms of functionality.

The Internet is a hierarchical computer network. At the top of hierarchy are relatively few routers containing very general information in routing tables. The tables do not contain information about every host on the Internet, but only about a few network destinations. Thus, high level pre-defined ontological information could be distributed in the same way among the top routers (Figure 1).

\(^7\) We use the term broker and router interchangeably.
Fig. 1. Conceptual illustration of a two-level distributed semantic pub/sub network. Top-level routers have only high level descriptions of ontologies from the lower level routers.

It is difficult to envision what this higher level information will be at this time, but we only need to take a look at Internet directories such as Google and Yahoo to get an idea of top level semantic knowledge. Both of these directories provide a user with only a few key entries as starting point for exploring the vast Internet information store. We see top level brokers exchanging only covering and advertisement information.

Lower in the Internet hierarchy routers maintain routing tables with destinations to specific hosts. Even though top level brokers use a common ontology, lower level brokers do not have to. For example, consider two different pairs of communicating applications: financial and medical. Financial applications are exchanging stock quotes, while medical are exchanging news about new drugs. These two application use different ontologies. The ontology information for each application can be distributed between multiple routers. These low level brokers will advertise more general descriptions of the ontologies they have to higher level brokers. Using this information, any new application will be able to locate the broker with specific ontologies. Any application wishing to integrate medical and financial information can create a mapping ontology between the financial and medical ontologies and provide a general description of the mapping ontology to higher level broker like in the previous case. We see that high level concepts can be used to route information between brokers who do not have access to specific ontologies. We can look at these general terms as very terse summaries of ontologies.

Our vision of a large scale semantic-based routing raises many questions:

- **top-level routing:** How to bridge multiple distributed ontologies to enable content routing? How can we avoid or reduce duplication of ontological
information among brokers? What is an appropriate high level generalization that can bring together different ontologies? How do semantic routing protocols look like?

– **lower-level routing:** How to efficiently store ontological information at routers? Large knowledge databases will probably require secondary storage beyond what is available at routers. How does this affect routing? If routers have to use covering at this level how can they dynamically control the generality of covering to affect network performance?

4 Related work

We are not aware of any previous work addressing the semantic routing problem in pub/sub systems. Most research on semantic has been done in the area of heterogeneous database integration. The main problem in this area is on enabling integration of heterogeneous information systems so that users can access multiple data sources in an uniform manner. One way of solving this problem is by using ontologies. Semantic information systems use an ontology to represent domain-specific knowledge and allow users to use the ontology terms to construct queries. The query execution engine accesses the ontology either directly or via an inference engine in order to optimize the query and generate an execution plan. Use of an ontology to generate an execution plan is central in determining the right source database and method for retrieving the required information. This allows uniform access to multiple heterogeneous information sources. The problem of adding semantic capability to pub/sub systems can be seen as an “inverse problem to the heterogeneous database integration problem. In semantic pub/sub systems, subscriptions are analogous to queries and events correspond to data, so now the problem is how to match data to queries.

Some systems \([8, 9]\) use inference engines to discover semantic relationships between data from ontology representations. Inference engines usually have specialized languages for expressing queries different from the language used to retrieve data, therefore user queries have to be either expressed in or translated into the language of the inference engine. The ontology is either global (i.e., domain independent) or domain-specific (i.e., only a single domain) ontology. Domain-specific ontologies are smaller and more commonly found than global ontologies because they are easier to specify. Additionally, there are systems that use mapping functions exclusively and do not have inference engines \([10, 11]\). In these systems, mapping functions serve the role of an inference engine.

Web service discovery is a process of matching user needs to provided services; user needs are analogous to events and provided services to subscriptions in a pub/sub system. Web service discovery systems \([12, 13]\) are functionally similar to a pub/sub system. During a discovery process, a web service advertises its capabilities in terms of its inputs and outputs. An ontology provides an association between related inputs or outputs of different web services. A user looks for a particular web service by searching for appropriate inputs and outputs according to the users needs. Relevant services are determined by either
exact match of inputs and outputs, or a compatible match according to ontology relationships.

The main push for using ontologies and semantic information as means of creating a more sophisticated application collaboration mechanisms has been from the Semantic Web community\(^8\). Recently their focus was on developing DAML+OIL—a language for expressing, storing and exchange of ontologies. Our vision of a distributed semantic pub/sub system is similar to that of the semantic web. The issues of distributing ontological information and bridging of different ontologies are common to both.

A system for distributed collaboration [14] creates a virtual network of proxies (functionally similar to brokers) using IP multicast connecting both data producers and consumers (users). Using common ontology, sources provide descriptions (metadata similar to subscriptions and events) of multimedia data they are providing and users provide descriptions of their capabilities. The metadata is distributed among proxies to create a semantic multicast graph along which data is distributed to interested users.

To improve scalability, peer-to-peer database systems are looking in the direction of semantic routing. HyperCuP [15] uses common ontology to dynamically cluster peers based on the data they contain. A cluster is identified using a more general concept than those associated with its members in the ontology. Concepts in the ontology map to cluster addresses so a node can determine appropriate route for a query by looking up more general concepts of the query terms in the concept hierarchy. Edutella [16] uses query hubs (functionally similar to brokers) to collect user metadata and present the peer-to-peer network as a virtual database which users query. All queries are routed though a query hub which forwards queries only to those nodes that can answer it.

5 Conclusions

In this paper we underline the limits of matching and content-based routing at syntactic level in pub/sub systems. We propose a solution for achieving semantic capabilities for local matching and look into the implications of using such a solution for content-based routing. We also present our vision on next-generation semantic-based routing. Our intent was to give rise to questions and ideas in order to improve existing content-based routing approaches and make them semantic-aware.

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