Comparative Analysis on Load Balancing Techniques in Cloud Computing

Amit Verma*, Sanjay Kumar Sharma and Bhavneet Kaur

1Department of Computer Science, Chandigarh Engineering College, Landran, Mohali – 140307, Punjab, India; dramitverma.cu@gmail.com, bhavneet14kaur@gmail.com
2Department of Computer Science, Chandigarh University, Gharuan (Mohali) - 160036, Punjab, India; talk4sanjaysharma@gmail.com

Abstract

Background/Objectives: Cloud computing is an arena that is ruling the world of information technology. Every user has its own definition for this technology as per their use. This paper is properly discussed document that describes the complete evolution of cloud computing from its beginning. Findings: With the presence of vast literature in field of load balancing, it was found confusion for the new scholars to find the startup point for their research in this field. Therefore, an exhaustive comparison has been made for the superior understanding of cloud evolution through various proposed algorithms from the past many decades, which will make the researchers possible to analyze the existing scenarios and a better way out to overcome the unsolved queries. Application/Improvements: The assessments between the algorithms will help the new researchers to analyze and opt for the parameters those need much more concentration to meet the required targets for better outcomes in the field.

Keywords: Cloud Computing, Information Extraction and Performance Measure, Load Balancing

1. Introduction

Computing, a term originated from the term compute which describes a way to attain the results from corresponding task a user had performed. Therefore, we can define computing as an attaining of results from evaluating the problem. Civilization of computer society is always on its peak in terms of the development. The most emerging technology of current era on which massive investigations and drastic progressions are coming is cloud computing. It is a drastic development in field of computing which has been emerged by witnessing the development that followed the pathway from distributing computing to parallel computing, followed by cluster computing and grid computing. Therefore, cloud computing is considered as a striking computing model which allows for the provisioning of resources on-demand, which in a layman language can be described as a virtual space in which users are allotted some space for storage which is accessible by the user ID and password that is actually their private space but along with this there is another kind of space, where data present on it has open access to people for its usage. The live example of who is Google.

Table 1 is presented that describes the tabular representation of the evolution in technology and network environment in computing. This paper we are going to study the algorithmic evolution of load balancing in computing.

| S No. | Era              | Technology     | Network Environment |
|-------|------------------|----------------|---------------------|
| 1     | From 1950        | ECL            | SNA                 |
| 2     | 1950 - 1970      | Custom Bipolar | SNA, DEC            |
| 3     | 1970 - 1980      | CMOS Micro     | TCP/IP and NFS      |
| 4     | 1980 - 1990      | CMOS Micro     | TCP/IP and NFS      |
| 5     | 1990 - till date | X86 Micro      | Netware             |

* Author for correspondence
2. Evolution Study via Algorithms

One of the fruitful ways to not only improve the throughput of systems but also enables the effective usage of resources and reduction in response time can be done by the load balancing. Many algorithms were designed to solve various issues related to computing but here we are going to discuss the evolution of the load balancing for various computing environments. Firstly, Table 2 presented here describes the comparisons between the few selected algorithms from different years of development.

Table 2. Comparison between existing scheduling algorithm

| Title     | Parameters | Findings | Environment | Tools     |
|-----------|------------|----------|-------------|-----------|
| LBDMC     | C_pattern, W_red, Th_red, L_red, D_emt | JDK      |
| MQCSS     | S_rate, T_m, C_i, M_sim, C_work, C_envi, C_sim |

DFD Level 0 is used to represent the basic overview of the system similarly, the DFD level 0 for evolution of technology describes that the information will be extracted from cloud by using various techniques like clustering, gridding, scheduling and many more.

Table 3 represents the DFD level 0 of technology evolution for load balancing in cloud computing.

Level 1 DFD for evolution of technology describes that to extract the information for load balancing in cloud would only be possible by using any sort of technique like clustering, scheduling, process migration, mark span, gridding on the data and along with the use of techniques the tool is required for the data extraction accordingly.

Table 4 shows the DFD Level 1 respectively.

DFD Level 2 is that level of data flow diagram that describes the detail prototype of complete system. Like as in Table 5 that displays the information extraction from cloud for load balancing. Here it has been described that the techniques used for load balancing are clustering, scheduling, process migration, mark span and gridding. And correspondingly the tools used on these techniques are ND tool, JDK, OpenVZ, CloudSim and GridSim respectively.

In author has presented an idea of scheduling parallel applications for multiuser, non-homogenous and large scale distributed systems. Target here is to design a system which has ability to handle idle cycles in network along with ability to handle clustered system and parallel processors. The algorithm proposed here is a better combination of three already existing algorithms that can handle multiple processors capabilities, various types of architecture and the variations in the processes. Here the decisions for scheduling will be done with the aim to reduce the turnaround time. At last the virtual processors for each application are gang scheduled to increase the efficiency of the system.

In author has presented an algorithm that runs with an aim to minimize response time and total time. These algorithms are used to develop an algorithm for query processing. These algorithms have result in more effectiveness cost and execution.

In author has presented a survey on distributed system design for load sharing. The paper is proposed with the source and server initiative approaches. Performance evaluation is done between the ten selected algorithms. Here a QFactor has been defined that ranked the various selected algorithms on the basis of their efficiency and fairness. After that the evaluation is done by using various mathematical and simulation techniques, it has been concluded that design decision is a critical issue and existing algorithms are providing effective solutions.

Table 6 describes the complete review of the papers from year 1900-2000. In author has developed a successful algorithm that can handle the load of the system through cluster nodes. When they switch on the cluster handles the node dynamically and according turn off to save the energy. Algorithm works at three level firstly at cluster based, secondly at OS level and lastly by application negotiation.

In author has discussed the issues related to the agent properties and load balancing. An algorithm of load balancing has been proposed that is based on communication in multiagent computing field. That has been implemented and correspondingly results have been evaluated. In algorithm, credit value is associated with every agent. And this credit value depends on the machine’s affinity, workload, mobility etc. Here the credit of every agent is investigated in the scenario of load imbalancement and then the lowest credit value agent will be drifted to the machine with the light load in system. Here experiment is held to analyze the comparison of improvement in both load balancing and its schemes oriented with workload.

In author has developed an idea of GHS i.e. Grid Harvest Service that gives dynamic and self-adaptive scheduling which can works in large applications in non-homogenous environment.
Table 3. DFD Level-0

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Comparative Analysis on Load Balancing Techniques in Cloud Computing

In author has proposed an algorithm that carries a capability to organize sensors present in WSN into clusters. Algorithms have ability to attain cluster head hierarchy in which energy saving increase with its each level.

In author has proposed a strategy for job grouping at runtime which in together have ability to perform simulation analysis. Quantity of jobs that can be processed in certain time can be determined by the processing of granularity size.

In author has developed an algorithm which is layered approach for load balancing that too in grid computing. Algorithm supports heterogeneity along with scalability. It also made concern for adaptability. The algorithm is on top is a tree like structure model.

In author has proposed a work for cloud workflow scheduling. It is made in concern to have pay as per use for execution and execution time. It was focused that algorithms must carry facilities for execution cost and time according to user’s input.

Table 4. DFD Level-1

Table 5. DFD Level-2

Table 6. Comparison between algorithms form 1900-2000

| Author        | Title | Year | Parameters | Algorithms Used | Pros | Cons | Tools / Tech used |
|---------------|-------|------|------------|-----------------|------|------|-------------------|
| S.A Khaled et al. [9] | GSHDS | 1900 | Flx E_l N ap N ap MTAT C comp A E_A E_l N ap N ap L_d S_ch N ap N ap | | | |
| Alan R. Hevner et al.[10] | QPDDS | 1979 | N ap N ap N ap G based N ap N ap C_e E_g E_e L_con | | | |
| Yung-Terng Wang et al. [11] | LSDS | 1985 | S_m N ap P_moc F_tol N ap N ap N ap N ap N ap C_e N ap N ap | | | |
| Author                | Title         | Year | Parameters | Algorithms Used | Pros | Cons | Tools / Tech used |
|-----------------------|---------------|------|------------|-----------------|------|------|-------------------|
| Eduardo Pinheiro et al. [12] | LBP_CBS      | 2001 | LBCBS      | P_wer E_ogy N_ap C_confi L_distrition N_ap | E_conserv. | P_conserv. N_ap P_et C_1 | ND-tool |
| Ka-Po Chow et al.[13]  | LBDMC         | 2002 | C_pattern W_load A_gnt H_v C_smart N_ap | CBS | Thhold L_red S_sp N_ap | P_loss | JDK N_ap |
| Ming Wu, et al. [14]  | SATS_H       | 2003 | A_1 P_mocc T_m M_min TA M_min GA S_adap TS T_reliant | P_loss C_time C_cost M_cost | GHS N_ap |
| Seema B. et al. [15]  | EHCAWN        | 2003 | E_ogy T_m N_ap EE Algo N_ap | L_ergy T_ergy N_ap C_cost | MSP BUA N_ap |
| Nithiapiday M et al. [16] | DL_GB_SA_GT    | 2005 | T_m N_ap N_ap J_A Sschedule A N_ap | Com_time P_vtime N_ap N_ap J_gg | G_Sim N_ap |
| Belabbas Yagoubi et al. [17] | DLBS          | 2006 | I_d C_v N_ap I_wde I_cluster I_gnd I_bal C_cont N_ap G_s O_graph C_1 | N_ap |
| Yun Yang et al. [18]  | SSTCCW        | 2008 | C_exec* T_m N_ap WSA N_ap | C_Mini C_exec N_ap N_ap R_App | SDW-C N_ap |
| Yi Zhao et al. [19]   | AD_LMA_VM     | 2009 | C_vendor T_m N_ap N_ap CB N_ap | F_Cost N_ap N_ap M_Jlow, A_A N | OpenVZ N_ap |
| Saeed Parsa et al.[20] | RASA          | 2009 | M_span N_ap N_ap M_span M_span RASA N_ap | R_MS N_ap N_ap D_task | C_exec G_Sim N_ap |
| Meng Xu et al. [21]   | MQCSS         | 2009 | SS_satel T_m M_span, C_v Schedule A T_cost N_ap | D_wrldfw C_exec T_exec Pople R_tabile | C_Sim N_ap |
| Mrs.S.Selvarani et al. [22] | ICBATS       | 2010 | P_mocc T_m N_ap T_gprod ABC I_schedule Comp_Cost N_ap N_ap SS_schedule, D_Factors | C_sim N_ap |
| Ke Liu et al. [23]    | CTCSAS-DC     | 2010 | C_t T_m N_ap N_ap CTC N_ap | C_exec* T_exec N_ap N_ap | C_Entirel SDW-C N_ap |
| Suraj Pandey et al. [24] | PSOHSW       | 2010 | R_metic T_m N_ap N_ap Schedule A PSO N_ap | C_Mini W_dis N_ap R_App AEC_2 | N_ap |
## Table 8. Comparison between algorithms from 2011-2015

| Author                     | Title            | Year | Parameter | Algorithms Used | Benefit (pro) | Cons | Tool used (tech) |
|----------------------------|------------------|------|-----------|-----------------|---------------|------|-----------------|
| T.Kokilavani et al. [25]   | LBMSMTS          | 2011 | $M_{\text{span}}$ | LBMM            | N             | N    | LBMM            |
| Ye et al. [26]             | SSWEC            | 2011 | $T_{\text{exec}}$ | SHEFT           | N             | N    | SHEFT           |
| Subasish Mohapatra et al. [27] | CHLBVM $C_L$   | 2013 | $T_m$      | ESCE            | N             | N    | ESCE            |
| P. Keerthika, et al. [28]  | HSLAB            | 2013 | $F_{\text{tol}}$ | MCSA            | N             | N    | MCSA            |
| Priyanka Gautam, et al. [29] | ERRLB            | 2014 | $C_r$, $T_m$ | $E_{\text{cpu}}$ | N             | N    | $E_{\text{cpu}}$ |
| Santanu Dam et al. [30]    | GAGEHL           | 2015 | $M_{\text{span}}$ | GHLBA           | N             | N    | GHLBA           |
| Salawu Nathaniel et al. [31] | DPAT$_d$         | 2015 | $C_r$      | DPAT-$dA$       | N             | N    | DPAT-$dA$       |
| M. Lavanya et al. [32]     | EL$_d$ S$_p$ PC  | 2015 | $T_t$      | ELB             | N             | N    | ELB             |
| Moon Suk Yeon et al. [33]  | MLL$_b$ HWB      | 2015 | $S_{cb}$   | PPAL$_d$        | N             | N    | PPAL$_d$        |
In\textsuperscript{19} author has proposed a Compare and Balance sampling based load balancing algorithm. In this paper virtual machine migration implementation has been worked on a simple and basic model that reduces the migration time of VM on the basis of storage. The results of above are successful.

In\textsuperscript{20} author has proposed an algorithm RASA with features of scalability and distribution. It is based on the study of traditional scheduling algorithms. At top of all, the proposed algorithm carries merits and overcome the demerits of these studied algorithms. Like it execute non-large tasks first as compare to others and support concurrent execution of large and small task in spite of execution of just large tasks. Results show low mark span as compare to early.

In\textsuperscript{21} author has presented a Multiple QoS Constrained Scheduling Strategy of Multi-Workflows. Here problems are handled by workflows that are related to requested services of different QoS requirements by multiple users at same time. This strategy works on the contribution of task sort algorithm and schedule algorithm. By this strategy the drawbacks of RASA dynamic workflow, issues related to execution cost and time had overcame.

In\textsuperscript{22} author has proposed an algorithm that is concerned with the job grouping costs. Aim in this paper is to task group scheduling it is because of job grouping optimization increases. Here algorithm measures cost of resources and computation of performance.

In\textsuperscript{23} has proposed an algorithm that is time cost algorithm for scheduling. This algorithm proves that cost of execution is decrease by around 15\% whereas execution time is decrease by 20\% within execution cost of user.

In\textsuperscript{24} author has described a heuristic which is based on PSO i.e. Particle Swam Optimization that will schedule cloud resources to applications. Here experiment has been conducted on application of workflow by timely varying its communication and computational cost and it was conclude that PSO is three times saver in comparison to existing BRS i.e. it has best resource selection and is nice distributor of workload.

Table 7 is presented that describes the complete review of the papers from year 2001-2010.

In\textsuperscript{25} author has studied that the basic min-min algorithm is the one that reduce the mark span as compare to others which somehow failed in load balancing scheduling but in this paper they proposed algorithm which on one hand decrease make span and in other hand increase the utilization of the resources. For this two phases are developed in first phase where basic min-min algorithm is executed in other hand non-utilized resources are made in use with maximum effectiveness.

In\textsuperscript{26} author has proposed a SHEFT algorithm that schedule workflow. It's an algorithm that elastically schedule workflow at runtime.

In\textsuperscript{27} author has discussed and presented the comparative study of various load balancing policies that are developed for effective and efficient service providing. But in our paper we have only used round robin algorithm for load balancing.

In\textsuperscript{28} author has proposed an algorithm for load balancing, architecture based user satisfaction, fault tolerance, availability and heterogeneity of resources. Along with this using Grid Sim tool kit has reduced the make span. And it results in better hit rate of make span and utilization of resources.

In\textsuperscript{29} author has observed that there a problematic issue of content storage which cause clustering and waiting queue for load balancing. Here the extended version of existing round robin algorithm has been proposed. This has resulted in reduced execution time and execution cost.

In\textsuperscript{30} author has discussed a strategy for load balancing to balance overwhelmed node using cloud analyst simulation tool. Then the results attained from this proposed algorithm are made compared to the existing algorithms.

In\textsuperscript{31} author proposed a dual-parameters adaptive timer decision algorithm for balancing the load and to simultaneously control the Ping-Pong handover in the network. And algorithm became successful to achieve 95\% level for load balancing.

In\textsuperscript{32} ELB algorithm has been proposed for scheduling of tasks to virtual machines that result in high throughput and low turnaround time. Here global queue has been implemented that results in increase of response time and efficiency.

In\textsuperscript{33} Packet processing algorithm has been proposed that deals with bottom level dynamic weight based agents. Those later results increase in response time.

Table 8 is presented that describes the complete review of the papers from year 2011-2015.

Some of the algorithms are selected from all algorithms of evolution. Whose tabular comparison is shown in Table 2 and the algorithms are discussed below:
Comparative Analysis on Load Balancing Techniques in Cloud Computing

Algorithm: Complete_Comet

Begin
  Initialize $C_L$ and sleep $T$
  For ($C_L=1$, $C_L<=L_{agent}$, $C_L++$) // where $C_L$ current location
    For each Location ($C_c=1$, $C_c<=C_{last}$, $C_c++$)
      Calculate the credits for each location
    
  If
    $C_{res}=I_{required}$ // Where $I_{required}$ is information required to be calculated for each load.
    Target meet
    Then
      Change $T_H$ // Where $T_H$ is the threshold.
    Exit(1)
  Else
    Migrate the agent with the smallest credit according to the selection and location policies
  End

Comet algorithm is proposed in\textsuperscript{13} that employs information collection and decision making policy. Here machines are synchronized and analyze load at periodic basis in contradiction with load threshold. This system has facility to report average. Here $T$ is a set of all tasks and $S$ is a set of available services. The major motive is to first submit the workflow of user with appropriate Quality of Services (QoS). After this the system will allocate the services for workflow and correspondingly will schedule the tasks according to QoS needs and environment load and variance at every phase of load balancing. After that selection decision is made accordingly. Corresponding algorithm is described above.

Input: Threshold load value $T_H$ (determined by profiling), information collection period $T$ (in seconds), number of agent’s $n$, hosts characteristics (address and other system information) and number of host’s $sp$.

In\textsuperscript{21} Multiple QoS Constrained Scheduling Strategy of Multi-Workflows has been introduced. Here problems are handled by the workflows that request the service of different QoS requirements by multiple users at same time. This strategy works on the contribution of task sort algorithm and schedule algorithm. By this strategy the drawbacks of RASA dynamic workflow, issues related to execution cost and time are overcame. Both the task sort algorithm\textsuperscript{21} and schedule algorithm\textsuperscript{21} are discussed above.

3. Conclusion

A complete evolution of computing from its preliminary stage to the current progressions has been discussed in this paper. Scheduling has been considered as an important issue for managing execution of applications in cloud. Therefore, progress that how load balancing in computing environment was done from early years to current year has been studied. Along with the assessments between the algorithms that have actually overcome the shortcomings of the existing algorithms, the parameters those were considered by various researchers which will help the new researchers to analyze and select the parameters those need much more concentration to meet the required functions are discussed in this paper.
Algorithm: Task Sort Algorithm

Begin
Initialize T, S

G_{GR}(t) // Where G_{GR} is a Get Ready Task
{ 
    while ( t is Ready )
    t=RT // Where RT is a ready task
}

ST(R_{g}, q) // Where ST is Sort Task
{ 
    while (t. RT)
    IT(t, q)
}

IT (t, q) // Where IT is Insert Task
{ 
    Insert t into q according to strategy
}
End

Algorithm: Schedule Algorithm

Begin
Initialize q, t

S_{Ch}(q,S) // Where S_{Ch} is a schedule
{ 
    while (q!=0)
    { 
        q= t // first task in queue
        s=gets(t, S)
        schedule t on s
        q= q – t
        S= S – r
    }
}

Get_{s}(t,S) // Where Get_{s} is a Get Service
{ 
    Select s.S
    (E_{time} and C_s) <= (T_{time} and T_{cost}) // Where E_{time}
    is Execution Time,
    C_s is cost, T_{time} is Total Time and
    T_{cost} is Total cost
    return s
}
End

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Appendix

- Process Migration, $P_{Mig}$ - Gridding, $G_{grid}$ - Query Processing in Distributed Database System, $QPPDS - Query Processing in Distributed Database System$, $GSHDS$ - Grid-based Grammar Based Algorithm, $C$ - Cost Effective, $E_{qps}$ - Effective query processing subsystem, $E_{exec}$ - Effective execution, $I_{concern}$ - Integral part concern, $C_{failure}$ - Component failure, $S_{task}$ - Scheduling, $C_{clustering}$ - Clustering. LSDD - Load Sharing in Distributed Systems, $S_{service}$ - Service, $P_{max}$ - Performance $F_{tol}$ - Fault Tolerance, $N_{app}$ - Not applicable, $C_{class}$ - Simple class considerations, $S_{sub}$ - Small subsets. LBP, CBS - Load Balancing and Unbalancing for Power and Performance in Cluster-Based Systems, LBCS - Load Balancing in Cloud Based System, $P_{wer}$ - Power, $E_{ngy}$ - Energy, $C_{conf}$ - Cluster configuration algorithm, $L_{dist}$ - load distribution algorithm, $E_{cons}$ - Energy conservation $P_{cons}$ - Power conservation, $P_{et}$ - Execution time prediction, $ND_{tool}$ - ND Tool, LBDMC - Load Balancing for Distributed Multiagent Computing, $C_{pattern}$ - Communication Pattern, $W_{load}$ - Workload, $A_{task}$ - Agents, $H_{hosts}$ - Hosts, $C_{comet}$ - Comet, $T_{hold}$ - Threshold, $L_{red}$ - Load Reduction, $S_{ch}$ - Scalability, $P_{loss}$ - Performance Loss, SATSHC - Self-adaptive Task Scheduling System for Non-dedicated Heterogeneous Computing, SATS - Self-adaptive Task Scheduling System, $A_{alloc}$ - Allocation, $P_{mnce}$ - Performance, $T_{time}$ - Time, $M_{min}$ - Mean-time task allocation algorithm, $M_{GA}$ - GA-Min-task group allocation algorithm, $S_{adapt}$ - Self-adaptive task scheduling algorithm, $T_{ref}$ - Task reallocation $P_{loss}$ - Performance loss, $C_{time}$ - Completion time, $C_{com}$ - Communication cost, $M_{cost}$ - Migration cost, GHS - Grid Harvest Services. EHCAWN - Energy Efficient Hierarchical Clustering Algorithm for wireless Sensor Network, $E_{ngy}$ - Energy, $T_{time}$ - Time, $E_{algo}$ - Energy Efficient Algorithm, $I_{engy}$ - Less energy consumption, $T_{cany}$ - Time complexity, $C_{cost}$ - Communication cost M$M$ - Medium
access protocol, BUA- Bottom- up approach. D,GB. \textit{FA}_{GT} Dynamic Job Grouping-Based Scheduling for Deploying Applications with Fine-Grained Tasks on Global Grid, T_m-Time, J\textit{g}_{group}J- Job Grouping Algorithm, S_{chdule-A} Schedule Algorithm, C_{\text{Sim}}-Communication time, P_{\text{overhead}}- Processing overhead, J_{\text{g}}-Job Grouping, C_{\text{Sim}- Grid Sim. D_{\text{LBSG}} Dynamic Load Balancing Strategy for Grid Computing DBLS Dynamic Load Balancing Strategy, L_d-Load, C_{\text{ct}}-Cost, I_{\text{site}}- intra-site algorithm, I_{\text{cluster}}- intra-cluster algorithm, I_{\text{grid}}- intra-grid algorithm, I_{\text{bal}}- Load balancing, C_{\text{com}}-Communication cost, G_{\text{MB}}- Grid simulators, O_{\text{re}} Real grid environment operations, C_{\text{CT}}- Clustering, AD_{LBA}MVM- Adaptive Distributed Load Balancing Algorithm based on Live Migration of Virtual Machines in Cloud, T_m-Time, C_{\text{nvrg}}-Convergence, C_{\text{Slow}}- Slow Migration for virtual machine AA_{\text{NSopt}} No Support for affinity and anti-affinity. RASA- RASA: A New Task Scheduling Algorithm in Grid Environment, M_{span}- Markspan, M_{m}\_m m- Min-min Algorithm, RASA-Resource Aware-Scheduling algorithm, R_{MS} reduce makespan, D_{dsk}- Tasks Deadlines, C_{exec}-Execution Cost, C_{\text{Sim}- Grid Sim. N_{ap}- Not applicable. MQCSS- A Multiple QoS Constrained Scheduling Strategy of Multiple Workflows for Cloud Computing, S_{\text{rate}}- Schedule Success rate, T_m-Time, C_{\text{ct}}-Cost, M_{span}- Markspan, T_{\text{Sort}}- Task Sort Algorithm, S_{chdule-A} Schedule Algorithm, D_{\text{widthb}}- Dynamic workflow, C_{exec}-Execution Cost, T_{exec}- Execution time, A_{\text{avail}}- Availability, R_{\text{Rel}}- Reliability, C_{\text{Sim}- Cloud Sim. ICBATS- Improved Cost-Based Algorithm For Task Scheduling in Cloud Computing, T_m-Time, P_{\text{min}}- Performance, ABC-ABC Algorithm, T_{\text{gpecd}}- Task grouping and scheduling algorithm, I_{\text{chdule}}- Improved Scheduling algorithm, C_{\text{Comp}}- Compilation Cost, S_{\text{chdule-A}}- Simultaneous Scheduling D_{\text{Factors}}- Dynamic Factors, C_{\text{Sim}- Cloud Sim. CTCASDC- A Compromised-Time-Cost Scheduling Algorithm in SineDevC for Instance-Intensive Cost-Constrained Workflows on Cloud Computing Platform, C_{\text{ct}}-Cost, T_m-Time, CTC-Compromised-Time-Cost Scheduling Algorithm, C_{exec}-Execution Cost, T_{exec}- Execution time, C_{\text{exec}}- Multiple cloud environment, SDW-C- Swinburne Decentralised Workflow for Cloud, PSOHSSW- A Particle, Swarm Optimization-based Heuristic for Swarming Workflows Applications in Cloud Computing Environment, R_{\text{simul}}- Resource Utilization, T_m-Time, S_{chdule-A} Schedule Algorithm, PSO- Particle Swarm Optimisation Algorithm, C_{\text{cost}}- Cost Minimum, W_{\text{dis}}- Workload Distribution, R_{\text{App}}- Real Applications, AEC_{\text{C}}- Amazon EC2, LBMSMTS- Load balanced Min-