Limits to catalysis in quantum thermodynamics

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Abstract

Quantum thermodynamics is a research field that aims at fleshing out the ultimate limits of thermodynamic processes in the deep quantum regime. A complete picture of thermodynamical processes naturally allows for auxiliary systems dubbed ‘catalysts’, i.e., any physical systems facilitating state transformations while remaining essentially intact in their state, like an auxiliary system, a clock, or an actual catalyst. In this work, we present a comprehensive analysis of the power and limitation of such thermal catalysis. Specifically, we provide a family of optimal catalysts that can be returned with minimal trace distance error after facilitating a state transformation process. To incorporate the genuine physical role of a catalyst, we identify very significant restrictions on arbitrary state transformations under dimension or mean energy bounds, using methods of convex relaxations. We discuss the implication of these findings on possible thermodynamic state transformations in the quantum regime.

1. Introduction

In chemical reactions, it is common that a certain reaction should in principle be allowed, but in reality cannot take place (or occurs at extremely low rates) because of the presence of some large energy barrier. Fortunately, the situation is sometimes redeemed by the presence of certain chemical substances, referred to as catalysts, which effectively lower the energy barrier across the transformation. That is to say, catalysts significantly increase the reaction rates. Importantly, these catalysts can remain unchanged after the occurrence of the reaction, and hence a small amount of catalytic substance could be used repeatedly and is sufficient to facilitate the chemical reaction of interest.

The basic principles of chemical reactions are governed by thermodynamic considerations such as the second law. There have specifically been a number of recent advances in the quest of understanding the fundamental laws of thermodynamics [1–6]. These efforts are especially focused on the quantum nano-regime, where finite size effects—either induced by system of interest, environment [7], or external fields used to govern time dependent interactions [8]—and quantum coherences are becoming increasingly relevant. One particularly insightful approach is to cast thermodynamics as a resource theory [2, 3, 9, 10], reminiscent of notions in entanglement theory [11–13]. In this framework, thermodynamics can be seen as the theory that describes conditions for state transformation $\rho_s \rightarrow \sigma_t$ from some quantum state to another under thermal operations (TO). The notion of TO means allowing for the full set of global unitaries which are energy preserving in the presence of some thermal bath. This is a healthy and fruitful standpoint, and allows the application of many concepts and powerful tools derived from information theory [14–16]. This class specifically entails all those thermodynamic processes that one can realistically physically implement in the presence of thermal baths, when encompassing real-world situations [17]. This seems a most natural view particularly in light of the emerging limitations\textsuperscript{5}.

\textsuperscript{5} Since Gibbs preserving maps in [1] induce the same pre-order structure in the state space as thermal operations for this classical regime (block-diagonal states), our results would apply to the paradigm of Gibbs preserving maps [28] as well.
In the context of thermal operations, catalysts emerge as ancillary systems that facilitate state transformation processes: there are cases where $\rho_\omega \rightarrow \sigma_\omega$ is not possible, but there exists a state $\omega_C$ such that $\omega_C \otimes \rho_\omega \rightarrow \omega_C \otimes \sigma_\omega$ is possible. The metaphor of catalysis is appropriate indeed: by using such a catalyst $\omega_C$, one is enabled to perform the thermodynamic transformation $\rho_\omega \rightarrow \sigma_\omega$, while returning the catalyst back in its exact original form. This is called exact catalysis. The inclusion of catalyst states in thermal operations serves as an important step in an eventual complete picture of quantum thermodynamics; it allows us to describe transformations in the full picture, where the system is interacting with an experimental apparatus, for example a clock system. The mathematical conditions for catalytic transformations have been studied [18, 19] in the context of majorization [20]. These works were extended to the thermodynamic setting in [1] by showing that one can obtain necessary and sufficient conditions for exact catalysis in terms of a whole family of generalized free energies [1]. The ordinary second law of ever-decreasing free energy is but a constraint on one of these free energies.

Naturally, for physically realistic scenarios inexact catalysis is anticipated, where the catalyst is returned except for a slight degradation. The loss of catalytic ability over time is often observed in chemical reactions, suggesting that catalytic substances often undergo slight changes in thermodynamic processes. In the quantum nano-regime, uncertainties such as in the initial state, imperfections in implementation of quantum operations, or fluctuations induced by quantum noise can serve to induce small changes in the catalyst. On physical grounds, it is unreasonable to suggest that the catalyst is returned in exactly the same way. Surprisingly, it has been shown [1] that in some cases, the conditions for catalytic transformations are highly non-robust against small errors induced in the catalyst. The form of the second law thus depends crucially on the measure used to quantify inexactness. In particular, if inexactness is defined in terms of small trace distance, then there is no second law at all: for any $\varepsilon > 0$, there exists a state $\omega_C$ such that for any two states $\rho_\omega$ and $\sigma_\omega$, starting from $\omega_C \otimes \rho_\omega$, one can get to some $\omega_C' \otimes \sigma_\omega$ via thermal operations, where $\omega_C'$ is $\varepsilon$-close in terms of trace distance to $\omega_C$. We refer to this effect as thermal embezzling: instead of merely catalyzing the reaction, energy/purity has possibly been extracted from the catalyst and used to facilitate thermodynamic transformations, while leaving the catalyst state arbitrarily close to being intact [21]. On physical grounds, such a setting seems implausible, even though it is formally legitimate. A clarification of this puzzle seems very much warranted.

A first hint towards a resolution may be provided by looking at how the error depends on the system size. The trace distance error $\varepsilon$ depends on the dimension of the catalyst states $\dim(\omega_C) = n$; nevertheless one can find examples of catalysts where $\varepsilon \rightarrow 0$ as $n$ approaches infinity. While examples show that in principle thermal embezzling may occur [1], hardly anything else is known otherwise. Indeed, it would be interesting to understand the crucial properties that distinguish between a catalyst and a non-thermal resource in thermodynamics. From a physical perspective, it seems highly desirable to understand to what extent the effect of embezzling can even occur for physically plausible systems.

In this work, we highlight both the power and limitations of thermal catalysis, by providing comprehensive answers to the questions raised above. Our first result is derived in the regime where both the Hamiltonians of the system and catalyst are trivial, in other words, proportional to the identity operator. This result concerns the analytical construction of universal catalyst states, which are able to facilitate any state transition on the system $S$ (with some fixed dimension $m$). We show that for a catalyst to be universal, it is equivalent to facilitating a specific state transition, intuitively speaking, the hardest possible transition on system $S$. By analyzing such a problem, we then construct a family of universal catalyst states depending on catalyst dimension $n$, that achieves the optimal trace distance error.

The second part of our results is derived for general Hamiltonians of the system and catalyst $H_S$, $H_C$. We identify two reasons where thermal embezzling cannot happen: (1) when the dimension of the catalyst is bounded, and (2) when the expectation value of energy of the catalyst state is finite. For both cases, we are able to derive non-zero bounds on the trace distance error, therefore showing that $\varepsilon$ cannot be arbitrarily small. These bounds were derived under the assumption that the catalyst states are diagonal in the energy eigenbasis. Case 2 is especially interesting, since it holds for catalyst Hamiltonians with unbounded energy eigenvalues, as long as the partition function $Z_C$ is finite. These results have been derived by making use of splitting techniques to simplify the optimization problems of interest. The techniques can also be used to obtain more specific results: not only can we prevent thermal embezzling (which makes a statement about the ability of a catalyst to facilitate any state transition), but given a pair of states $\rho_\omega$, $\sigma_\omega$, one can obtain state-dependent lower bounds on the trace distance error as well.

2. Results

2.1. The power of thermal embezzling

We begin by exploring the case for trivial Hamiltonians, where it is known that thermal embezzling can occur. This is also the simplest case of thermodynamics in resource theory [1], when all energy levels are fully...
degenerate, and the Hamiltonian is simply proportional to the identity operator. In this regime, thermal states are simply maximally mixed states, and all unitary operations are allowed, a setting otherwise known as noisy operations. Entropy and information, instead of energy, become the main quantity that measures the usefulness of resources. In such cases, the sole conditions governing a transition from some quantum state \( \rho \) to \( \sigma \) is that the eigenvalue vector of \( \rho \) majorizes that of \( \sigma \) \[2\]. Majorization is commonly denoted as \( \rho \succ \sigma \); it also implies that entropy can never decrease under noisy operations \[1\].

To investigate thermal embezzling in this setting, one asks if given fixed \( m, n \), what is the smallest \( \varepsilon \) such that there exists a catalyst state \( \omega_C \) that satisfies
\[
\omega_C \otimes \frac{1}{m} \succ \omega_C' \otimes \ket{0} \bra{0},
\]
where the trace distance \( d(\omega_C, \omega_C') \) between the initial catalyst \( \omega_C \) and final catalyst \( \omega_C' \) is not greater than \( \varepsilon \). This trace distance is used as a measure of catalytic error throughout our analysis. If some catalyst pair \((\omega_C, \omega_C')\) satisfies the condition in equation (1) with trace distance \( \varepsilon \), then it also facilitates \( \omega_C \otimes \rho \rightarrow \omega_C' \otimes \sigma \) for any \( m \)-dimensional states \( \rho, \sigma \). This is because a pure state majorizes any other state, while the maximally mixed state \( 1/m \) is majorized by any other state.

Since majorization conditions depend solely on the eigenvalues of the density matrices \( \omega_C \) and \( \omega_C' \), one can phrase this problem of state transformation in terms of a linear minimization program over the catalyst states, diagonal and ordered in the same basis (see appendix). In fact, the eigenvalues of \( \omega_C \), \( \omega_C' \), which give rise to the optimal trace distance error can be solved by such a linear program, although for general values of \( n \) and \( m \), these eigenvalues are non-unique, and it is harder to construct an analytical solution. Whenever \( m \geq 2 \) and \( n = m^a \) where \( a \geq 1 \) is an integer, we provide an analytic construction of catalyst states, which we later show to be optimal for the state transformation in equation (1). Let the initial catalyst state be \( \omega_C = \sum_{i=1}^{n} \omega_{i} \ket{i} \bra{i} \), where \( \omega_{i} = m/(1 + (m - 1) a) \), and
\[
\omega_{i} = \begin{cases} \omega_{i} m^{-\log_{m} a} & \text{if } 2 \leq i \leq n/m, \\ 0 & \text{if } i > n/m. \end{cases}
\]
Note that our catalyst state \( \omega_C \) does not have full rank, and this is crucial for the majorization condition in equation (1) to hold, since \( \rho \succ \sigma \) implies that rank(\( \rho \)) \( \leq \) rank(\( \sigma \)), and the joint state \( \omega_C \otimes \ket{0} \bra{0} \) can have at most rank \( n \). The final state of the catalyst \( \omega_C' \) can be obtained from \( \omega_C \) by subtracting a small value \( \varepsilon \) from the largest eigenvalue \( \omega_{i} \) and distributing the amount \( \varepsilon \) equally over the indices \( i > n/m \). This causes \( \omega_C' \) to be a state of full rank \( n \). We show that this family achieves trace distance error
\[
d_{m,n} = \frac{m - 1}{1 + (m - 1) \log_{m} n},
\]
which we prove by mathematical induction to be optimal, given fixed \( m, n \) where \( n = m^a \) (see section B.1 of the appendix). The scenario \( n = m^a \) can be seen as follows: if the system is a particle in an \( m \)-dimensional Hilbert space, the catalyst consists of \( a \) number of such particles. The optimal error as presented in equation (3) scales with the number of particles \( a \) in the catalyst.

Figure 1 compares our final catalyst state with the state
\[
\hat{\omega}_C = \frac{1}{C(n)} \sum_{j=1}^{n} \ket{j} \bra{j},
\]
with \( C(n) = \sum_{i=1}^{n} 1/j \) being the normalization constant. The family \( \hat{\omega}_C \) was proposed in \[21\] for embezzling in the LOCC setting. In figure 2, we compare the trace distance error achieved by catalyst \( \hat{\omega}_C \) from \[21\] with the error achieved by our catalyst \( \omega_C \). We see that for small dimensions, our catalyst outperforms \( \hat{\omega}_C \), however asymptotically the error scales with \( \log n \) for both catalysts.

### 2.2. The limits of thermal embezzling
In this section, we are interested in finding additional physical restrictions which prevent thermal embezzling. To do so, we look at general Hamiltonians \( H_S, H_C \) of both the system and catalyst, where the energy of the system comes into play. The total Hamiltonian is simply \( H = H_S + H_C \), without any interaction terms. In \[1\], it is shown that the monotonicity of quantum Rényi divergences \[22\] (for \( \alpha > 0 \)) form the necessary conditions for state transformations. More precisely, for arbitrary \( \rho_S \) and \( \rho'_S \), if \( \rho_S \rightarrow \rho'_S \) is possible via catalytic thermal operations, then for all \( \alpha > 0 \),
\[
D_{\alpha}(\rho_S \parallel \tau_S) \geq D_{\alpha}(\rho'_S \parallel \tau_S)
\]
holds, where \( \tau_S \) is the thermal state of system \( S \), at temperature \( T \) of the thermal bath.
Equation (5) implies that one can use the monotonicity of Rényi divergences to find lower bounds on thermal embezzling error for state transformation between arbitrary states $\rho$ and $\rho'$. For simplicity, we present the case where $\rho$ and $\rho'$ are diagonal (in the energy eigenbasis of $H_S$). The case for arbitrary states can be treated similarly, and details are given in appendices C.2 and D.2.

For the case where two states $\rho$ and $\sigma$ are diagonal, the Rényi divergences are defined as

$$D_{\alpha}(\rho \parallel \sigma) := \frac{1}{\alpha - 1} \log \sum_i \rho_i^\alpha \sigma_i^{1-\alpha},$$

where $\{\rho_i\}$, $\{\sigma_i\}$ are the eigenvalues of $\rho$ and $\sigma$.

Again, for states $\rho_S$ and $\rho'_S$ diagonal, it suffices to look at a single transformation

$$\omega_C \otimes \tau_S \rightarrow \omega'_{C} \otimes \Pi_{\text{max}}^S,$$

where $\Pi_{\text{max}}^S = |E_{\text{max}}^S\rangle\langle E_{\text{max}}^S|$ is the pure energy eigenstate with energy $E_{\text{max}}^S$, the maximum energy eigenvalue of system $S$. Note that both $\tau_S$ and $\Pi_{\text{max}}^S$ are diagonal in the energy eigenbasis. Similar to the case of trivial Hamiltonians, the process in equation (7) is a sufficient condition for arbitrary diagonal state conversions. More precisely, if equation (7) holds, then for any $\rho_S$ and $\rho'_S$ diagonal in the same energy eigenbasis, $\omega_C \otimes \rho_S \rightarrow \omega'_{C} \otimes \rho'_S$ is also possible. This is stated formally and proven in lemma A.3, found in the appendix.

We also assume $\omega_C$ and $\omega'_{C}$ to be diagonal in the energy eigenbasis of $H_C$ [1]. This can be written as the following minimization problem

$$\epsilon := \min \frac{1}{2} \| \omega_C - \omega'_{C} \|_1$$

s. t. $\forall \alpha \geq 0, D\left( \omega_C \otimes \tau_S \parallel \tau_{CS} \right) \geq D\left( \omega'_{C} \otimes \Pi_{\text{max}}^S \parallel \tau_{CS} \right), \quad 0 \leq \omega_C, \omega'_{C} \leq 1,$$

where $\tau_{CS} = \tau_C \otimes \tau_S$ is the thermal state of the catalyst and system. The system Hamiltonian $H_S$ is assumed to be finite.
A first step towards solving equation (8) is to relax the problem to consider only a fixed \( \alpha \).

\[
\epsilon_\alpha := \min \frac{1}{2} \| \alpha_{\rho} - \alpha_{\sigma} \|_1
\]

s. t. \( D_H \left( \alpha_{\rho} \otimes \rho_{\tau} \| \tau_{CS} \right) \geq D_H \left( \alpha_{\sigma} \otimes \Pi_{\max}^E \| \tau_{CS} \right), \quad 0 \leq \alpha_{\rho}, \alpha_{\sigma} \leq 1. \tag{9}
\]

We know that any \((\alpha_{\rho}, \alpha_{\sigma})\) feasible for equation (8) is also feasible for equation (9). Therefore, for any \( \alpha \geq 0 \), \( \epsilon \geq \epsilon_\alpha \). By choosing \( \alpha \) one can arrive at much simpler optimization problems that provide lower bounds for the trace distance error. We apply this to study two cases, detailed as below.

1. **Bounded dimension:** Consider the case where both the system \( H_S \) and catalyst Hamiltonians \( H_C \) have fixed dimensions, and denote the maximum energy eigenvalues as \( E_{\max}^S, E_{\max}^C \) respectively. By choosing \( \alpha \to \infty \), we can obtain a lower bound for \( \epsilon \) in equation (8) based on these parameters. For \( \alpha \to \infty \), \( D_\infty (\rho \| \sigma) = \log \frac{P_i}{q_i} \), where \( P_i, q_i \) are eigenvalues of the states \( \rho, \sigma \) respectively. Recall that we have assumed that \( \alpha_{\rho} \) and \( \alpha_{\sigma} \) are diagonal in the same basis, which we take to be the energy eigenbasis. Equation (9) can be rewritten as

\[
\epsilon_\infty = \min \frac{1}{2} \sum_i |\alpha_i - \alpha'_i| \quad \text{s. t. } \max_i \frac{Z_S}{e^{-\beta E_{\max}^S}} \max_i \frac{\alpha_{i}'}{\tau_i}, \quad \alpha_i, \alpha_i' \geq 0 \quad \forall \ i, \quad \sum_i \alpha_i = \sum_i \alpha_i' = 1, \tag{10}
\]

where \( \tau_i = Z^{-1} \alpha_i e^{-\beta E_i} \) are the probabilities defined by the thermal state of the catalyst Hamiltonian, and \( Z_S, Z_C \) are partition functions of the system and catalyst respectively. To solve this problem, we note that the optimal strategy to maximize the quantity \( \max_i \alpha_i/\tau_i \) within the \( \epsilon \) ball of \( \alpha_{\rho} \) is to increase one of the eigenvalues by \( \epsilon \), so that the quantity \( \max_i (\alpha_i + \epsilon)/\tau_i \) is maximized. With further details in the appendix, we show that the trace distance error can therefore be lower bounded by

\[
d_{\text{opt}} \left( H_S, H_C \right) = \epsilon \geq \left( \frac{Z_S}{e^{-\beta E_{\max}^S}} - 1 \right) \frac{\sum_i \sqrt{P_i} q_i}{Z_C}. \tag{11}
\]

The bound in equation (11) depends on both the minimal population of the thermal state for system and catalyst. Although this bound is valid for arbitrary finite-dimensional Hamiltonians, it is not tight. Indeed, in the case of trivial Hamiltonians where all states have constant energy value, normalized to 0, the partition functions \( Z_S, Z_C \) reduce to the dimension \( m, n \) of the system and catalyst. This bound then yields \( d_{\text{opt}} (0_S, 0_C) \geq (m - 1)/n \), which is much weaker than the optimal trace distance (which scales with \( \log n \)) that we derived in equation (3).

2. **Hamiltonians with unbounded energy levels:** A more general result holds for unbounded dimension and energy levels where the partition function \( Z_C \) is finite. More precisely, for such cases, we show that setting an upper bound on the average energy of the catalyst state limits thermal embezzling.

Let us now explain the proof of our results. Consider some catalyst Hamiltonian \( H_C \) with unbounded energy levels \( [E_{\max}^C] \). For simplicity, we restrict ourselves to the case where the catalyst states are diagonal in the energy eigenbasis, and assume the system Hamiltonian to be trivial with dimension \( m = 2 \). The resulting bound is found in equations (19) and (22). A more general derivation involving arbitrary system Hamiltonians may be found in the appendix.

A) **Formulation of the problem:** Consider the minimization of catalytic error under the relaxed constraint that monotonicity for the \( \alpha \)-Rényi divergence is satisfied. Using equation (9) with \( \alpha = 1/2 \), for diagonal states \( \rho, \sigma \) with eigenvalues \( p_i, q_i \), equation (6) can be evaluated as \( D_{1/2} (\rho \| \sigma) = -2 \log \sum_i \sqrt{p_i} q_i \). By substituting \( H_S = 0_S \), the first constraint can be simplified as follows

\[
D_{1/2} \left( \alpha_{\rho} \otimes \frac{\rho}{2} \| \tau_C \otimes \frac{\tau}{2} \right) \geq D_{1/2} \left( \alpha_{\sigma} \otimes |0\rangle \langle 0| \| \tau_C \otimes \frac{\tau}{2} \right). \tag{12}
\]

\[
D_{1/2} \left( \alpha_{\rho} \| \tau_C \right) \geq D_{1/2} \left( \alpha_{\sigma} \| \tau_C \right) + D_{1/2} \left( |0\rangle \langle 0| \| \frac{1}{2} \right) \tag{13}
\]

\[
\sqrt{\sum \alpha_i^{1/2} e^{-\beta E_i^{1/2}}} \leq \sum \alpha_i^{1/2} e^{-\beta E_i^{1/2}}. \tag{14}
\]

Equation (14) follows by the additivity of all Rényi divergences, and equation (14) is obtained by evaluating all \( D_{1/2} \) terms. Furthermore, we want that the initial catalyst state to have an expectation value of energy no larger than some finite \( E \). In summary, we now look at the minimization of trace distance under the following constraints.
\[ \epsilon := \min \frac{1}{2} \sum_i |\omega_i - \omega'_i| \]
\[ \text{s. t. } \sum_i \sqrt{|\omega_i \gamma^{E_i^C}|} \geq \sum_i \sqrt{|\omega_i \gamma^{E_i^C}|}, \quad \omega_i, \omega'_i \geq 0 \quad \forall \ i, \quad \sum_i E_i^C \omega_i \leq E, \tag{15} \]

where \( \gamma = e^{-\beta/2} \in (0, 1) \). As such, this is an intricate problem, as it is a non-convex problem both in \( \omega_i \) and \( \omega'_i \).

In the subsequent steps, our goal is to show that \( \epsilon \) is lower bounded by a non-zero constant, by making use of the techniques of convex relaxations of optimization problems.

**B) Splitting a relaxed minimization problem:**

The key idea to proceed is to suitably split the problem into two independent optimization problems in a relaxation, which can be individually assessed. The starting point of this approach is rooted in the observation that for any \( \omega_i, \omega'_i \in [0, 1] \), the following inequality holds true,

\[ \omega_i^{3/2} - 2\omega_i^{1/2} \leq \left| \omega_i - \omega'_i \right|^{1/2} - \omega'_i/3. \tag{16} \]

Since requiring the rhs of equation (16) to be positive is less stringent compared to the lhs, one can now further use it to obtain a lower bound for the minimization in equation (15). By defining a new variable \( x_i = |\omega_i - \omega'_i| \), we can define a new minimization problem

\[ \zeta := \min \frac{1}{2} \sum_i x_i \]
\[ \text{s. t. } \sum_i \sqrt{|\omega_i \gamma^{E_i^C}|} \geq \frac{1}{3} \sum_i \sqrt{|\omega_i \gamma^{E_i^C}|}, \quad x_i, \omega_i \geq 0 \quad \forall \ i, \quad \sum_i E_i^C \omega_i \leq E, \tag{17} \]

and note that \( \epsilon \geq \zeta \). One can see now that the variables \( x_i, \omega_i \) are independent from each other. This allows us to first perform a minimization of the function \( \sum_i \sqrt{|\omega_i \gamma^{E_i^C}|} \) for constraints involving \( \omega_i \) only.

**C) Invoking energy constraints to provide lower bound:**

The energy constraint on \( C^{E} \) plays a crucial role in lower bounding \( \zeta \). Intuitively, when such a constraint is placed for some finite \( E \), it implies that the probability of populating some relatively low energy levels cannot be vanishingly small. We prove this with more rigour in the appendix. Along this line of reasoning, one concludes that for the minimization

\[ \epsilon_1 := \min \sum_i \sqrt{|\omega_i \gamma^{E_i^C}|} \]
\[ \text{s. t. } \omega_i \geq 0 \quad \forall \ i, \quad \sum_i E_i^C \omega_i \leq E, \tag{18} \]

\( \epsilon_1 > 0 \) has to be strictly positive. More precisely,

\[ \epsilon_1 = \max_{W \in [0,1)} \frac{W^{E^C_{\min}}}{E/j(W)}, \tag{19} \]

where \( j(W) = \min \{j: E_{j+1}^C > E/(1 - W)\} \). A derivation of this expression can be found in the appendix.

**D) Merging both problems:**

After obtaining a lower bound for the problem in equation (18), we recombine the two problems into equation (17) to obtain

\[ \min \frac{1}{2} \sum_i x_i \]
\[ \text{s. t. } \sum_i \sqrt{|\omega_i \gamma^{E_i^C}|} \geq \frac{1}{3} \epsilon_1, \quad x_i \geq 0 \quad \forall \ i. \tag{20} \]

This is a quadratic optimization problem in the variables \( \sqrt{|\omega_i \gamma^{E_i^C}|} \), hence it is easy to obtain the Lagrange dual of this problem, which takes on a very simple form

\[ \min - \frac{1}{4} \lambda \sum_i \gamma^{2E_i^C} + \lambda \epsilon_1 \quad \text{s. t. } \lambda \geq 0, \tag{21} \]

involving the simple minimization of a quadratic function w.r.t. \( \lambda \). Solving this we arrive at a lower bound

\[ \epsilon \geq \zeta \geq \frac{1}{2} \cdot \frac{1}{9} \epsilon_1^2 > 0, \tag{22} \]

where \( Z_C = \sum_j \gamma^{2E_j^C} = \sum_j e^{-\beta E_j^C} \) is the partition function of \( HC \). We summarize our findings in table 1.

### 3. Discussion and conclusion

In summary, we have carefully investigated the power and limitations of thermal embezzling under different physical scenarios, a setting that should be taken into account when considering the ultimate limits of...
thermodynamic transformations and work extraction under any physical mechanism. While in the fully degenerate Hamiltonian case, we have seen that the effect can be powerful, under physically ubiquitously common settings, it is very much limited. Based on very physical considerations, i.e. that catalysts have Hamiltonians, we resolve the puzzle of thermal embezzling, for all catalysts diagonal in the energy eigenbasis. In this way, we make a significant contribution to the quest for achieving a complete understanding of thermodynamic laws in the quantum world.

The bounds on dimensionality are closely related to energy restrictions. While placing an upper bound on the dimension directly implies an upper bound on the average energy, the reverse statement is not generally true. However, if one restricts not only the expectation value of the energy distribution, but also its variance to be finite, then this is almost equivalent to placing a dimension restriction. For example, given any non-degenerate Hamiltonian $H_C$ with unbounded eigenvalues, consider the set of catalyst states such that the average energy and variance of a given catalyst is finite. Then by the Chebyshev inequality one can understand that this is equivalent to introducing a cut-off on the maximum energy eigenvalue (and therefore on the dimension). We note that it is easy to see that e.g. for the harmonic oscillator the variance is not always bounded whenever the mean energy is bounded.

In the case of infinite-dimensional Hamiltonians, we have also shown that for certain classes of catalyst Hamiltonians, explicit bounds can be derived on the trace distance error of a catalyst when the average energy is finite. Our results have covered a large range of Hamiltonians which are commonly found in physical systems, including the important case of the Harmonic oscillator in free systems, with the minimal assumption that partition function $Z_C$ is finite, which holds for all systems for which the canonical ensemble is well-defined. However, we know that thermal embezzling can be arbitrarily accurate as the dimension grows, at least in the simplest case of the trivial Hamiltonian. This implies that there will be specific cases of infinite-dimensional Hamiltonians where simple bounds on average energy do not give explicit bounds on the thermal embezzling error. We suspect that this may be true for systems with unbounded dimension, but bounded Hamiltonians. The reason is that if dimension is unbounded, then there must exist an accumulation point in the energy spectrum. The subspace of this accumulation point will be very similar to the trivial Hamiltonian.

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### Appendix

In these appendices we fully elaborate our findings on thermal catalysis. We begin in section A.1 by explaining the similarities and subtle differences between thermal embezzling and embezzling in the LOCC setting. The Rényi divergences and their relation to thermal operations are detailed in section A.2. Proceeding to section B, we focus on thermal embezzling for trivial Hamiltonians with fixed dimensions. On the one hand, we investigate the problem of finding a catalyst which allows us to perform thermal embezzling with minimum possible error in trace distance. We detail the proofs on our construction of a catalyst family (given dimension parameters for both system of interest and catalyst), and prove that our construction achieves the optimal embezzling error.

On the other hand, by placing restrictions on the dimension, we derive non-zero lower bounds for the embezzling error, considering the arbitrary system and catalyst Hamiltonians. The proofs are detailed in section C. Some technical background on $\alpha$– Rényi divergences and their relation to thermodynamic operations are given. Lastly, in section D we focus on infinite-dimensional Hamiltonians, with unbounded
Appendix A. Embezzling and catalysis

A.1. Thermodynamics as a resource theory

Resource theories are frameworks useful in identifying states which are valuable, under specific classes of allowed operations and states given for free. A state is a valuable resource if one can use it to create many other states under the set of allowed operations. Thermodynamics can be viewed as a resource theory [2, 3], where the allowed operations are the so-called thermal operations. They are summarized as follows: considering a system $S$, given a state $\rho_S$ and the Hamiltonian $H_S$, one can

1. For any bath system $B$ with Hamiltonian $H_B$, attach any thermal state $\tau_B = e^{-\beta H_B}/\text{tr}[e^{-\beta H_B}]$ to $\rho_S$, where $\beta = 1/kT$ is the inverse temperature,

2. Perform any unitary $U$ over the global system $SB$ that commutes with the total Hamiltonian, i.e. $[U, H_S + H_B] = 0$,

3. Trace out the bath $B$.

Recently, the framework of thermal operations was used to prove a second law [1] by including catalytic effects. This is because there exist certain states $\rho$ and $\sigma$ such that via thermal operations $\rho \approx \sigma$, but $\rho \otimes \omega_C \rightarrow \sigma \otimes \omega_C$ for some state $\omega_C$. More precisely, catalytic effects can be accounted for by adding a fourth rule, i.e.

4. for any catalyst system $C$ with Hamiltonian $H_C$, attach any additional catalyst state $\omega_C$, as long as the returned state $\omega_C$ is $\varepsilon$-close to its original state $\omega_C$.

to the set of allowed operations. One can now ask, given $\rho_S$, what are the states $\rho_S'$ such that $\rho_S \rightarrow \rho_S'$ is possible via inexact catalytic thermal operations? More precisely, do there exist $\omega_C$, $\omega_C'$ which are $\varepsilon$-close to each other, such that $\omega_C \otimes \rho_S \rightarrow \omega_C' \otimes \rho_S'$?

Depending on $\varepsilon$ and the measure of closeness used, the conditions for $\rho_S \rightarrow \rho_S'$ to occur can vary. For example, if $\varepsilon$ is required to be zero, i.e. the catalyst must be returned in its exact form, then [1] shows for any $\rho_S$ and $\rho_S'$ such that $\rho_S \rightarrow \rho_S'$ is possible via catalytic thermal operations, a whole set of Rényi divergences must necessarily decrease. In the next section A.2, we define the Rényi divergences and state the results of [1] in detail. On the other hand, if $\varepsilon$ is measured in terms of trace distance between the initial and final catalyst only, [1] also proves that for any $\varepsilon > 0$, the state transformation conditions are trivial, i.e. any $\rho_S$ can be transformed to any $\rho_S'$. We denote thermal embezzling as the phenomenon wherein by requiring only the initial and final catalyst to be close in terms of trace distance, one can achieve $\rho_S \rightarrow \rho_S'$ for any $\rho_S$, $\rho_S'$.

Another well-studied example of a resource theory is entanglement theory, where the allowed operations are those that can be implemented using local operations and classical communiaction (LOCC), while free states are the set of separable states. The interconversion of resource states in entanglement theory has been studied intensively, and has also provided insight into the resource theory of thermodynamics.

Embezzling states were originally introduced for the LOCC setting in [21]. An entangled state $|\nu(n)\rangle_{AB} \in C^n \otimes C^n$ shared between two parties $A$ and $B$ can be used as a resource to prepare some other state (of much smaller dimension),

$$|\nu(n)\rangle_{AB} \rightarrow_{1-\varepsilon} |\psi\rangle_{A'B'}, \quad (A.1)$$

where $\dim(A^nB^n) \ll \dim(AB)$ and $\dim(A^nB^n) \approx \dim(AB)$. The fidelity between the actual final state with $|\nu(n)\rangle_{A'B'}$ is denoted by $1 - \varepsilon$, such that $\varepsilon$ goes to zero when $n$ goes to infinity. This enables the approximate preparation of the state $|\psi\rangle_{AB}$, while the embezzling resource state is also left close to its original state. Such a preparation can even be achieved simply via local operations (LO). The family $|\nu(n)\rangle_{AB}$ is called a universal embezzling state if it enables the preparation of any $|\psi\rangle_{A'B'}$. While this seemingly violates entanglement monotonicity under LOCC operations, one quickly realizes that it is really because the closeness in entanglement content of $|\nu(n)\rangle_{AB}$, $|\nu(n)\rangle_{A'B'}$ depend not only on the fidelity, but also the dimension. Hence entanglement is exhausted to prepare $|\psi\rangle_{A'B'}$, while $|\nu\rangle_{A'B'}$ remains close to intact on the whole. However, there is also something special about embezzling states, in the sense that a maximally entangled state does not serve as a good embezzler. In [13], a comprehensive study about the general characteristics of embezzling states was conducted, providing insight into the necessary structure of a state to be a good embezzler. The power of embezzling in LOCC has been applied in several areas of quantum information, such as coherent state exchange protocols [12], projection games [23], or as a theoretical tool in proving the quantum reverse Shannon theorem [24].
There are some similarities between thermal embezzling and LOCC embezzling; however, also many distinctive features exist. Most significantly, in thermodynamic systems, the Hamiltonian which determines the evolution of the system plays an important role in state conversion conditions [25]. This feature is absent in LOCC embezzling. We summarize the similarities and differences of LOCC and thermal embezzling in Table A1.

### A.2. Rényi divergences as thermal monotones

In this section we detail the conditions for state transformation under catalytic thermal operations, which are closely related to the Rényi divergences. The simplest case of catalytic thermal operations is when all Hamiltonians $H_B, H_C$ are trivial. For arbitrary states $\rho$ and $\sigma$, $\rho \to \sigma$ is possible if and only if $\rho \succ \sigma$ [2]. In the case where $H_B$ or $H_C$ are generally non–trivial, state conversion conditions are affected by the involved Hamiltonians. More precisely, instead of majorization, we need to consider the monotonicity of Rényi divergences as a (necessary) condition for state transformations. These conditions are used later in sections C and D to investigate the limits of thermal embezzling. Let us first define these quantities in definition A.1.

**Definition A.1 (Rényi divergences [22])**. Given arbitrary states $\rho$, $\sigma \succ 0$, for $\alpha \in [0, \infty)$, the Rényi divergence of $\rho$ relative to $\sigma$ is defined as

$$D_\alpha (\rho \parallel \sigma) := \frac{1}{\alpha - 1} \log \left[ \text{tr} \left( \frac{1-a}{\alpha} \rho^\alpha \frac{1-a}{\alpha} \sigma^\alpha \right)^a \right]$$

(A.2)

For $\rho$, $\sigma$ diagonal in the same basis, let $p = (p_1, p_2, ..., p_n)$ and $q = (q_1, q_2, ..., q_n)$ denote the eigenvalue vectors of the $\rho$, $\sigma$ respectively. Then the Rényi divergences reduce to the form

$$D_\alpha (\rho \parallel \sigma) = D_\alpha (p \parallel q) := \frac{1}{\alpha - 1} \log \sum_i p_i^\alpha q_i^{1-\alpha}.$$  

(A.3)

It has been shown that for diagonal states $\rho$, the quantities $D_\alpha (\rho \parallel \tau)$ are thermal monotones for all $\alpha \geq 0$, where $\tau$ is the thermal state of the system of interest. For arbitrary quantum states, $D_\alpha (\rho \parallel \tau)$ are thermal monotones for $\alpha \geq 1/2$ as well. Intuitively, this implies that thermal operations can only bring the system of interest closer to its thermal state with the same temperature $T$ as the bath [1]. We detail this in lemma A.2.

**Lemma A.2 (Monotonicity under thermal operations [1])**. Given some Hamiltonian $H_A$, consider arbitrary states $\rho_A, \rho'_A$ where $\rho_A \to \rho'_A$ is possible via catalytic thermal operations. Denote by $\tau_A$ the thermal state of system $A$. Then for any $\alpha \in [1/2, \infty)$,

$$D_\alpha (\rho_A \parallel \tau_A) \geq D_\alpha (\rho'_A \parallel \tau_A).$$

(A.4)

Furthermore, for any $\rho_A, \rho'_A$ diagonal in $H_A$, equation (A.4) holds for all $\alpha \geq 0$ if and only if $\rho_A \to \rho'_A$ is possible via catalytic thermal operations.

In essence, lemma A.2 implies that the monotonicity of Rényi divergences are necessary conditions for arbitrary state transformation, and for the case of states diagonal (in the energy eigenbasis), they are also sufficient. Let us also use a notation which was introduced in [25] for diagonal states: we say that there exists a catalyst $\omega$ such that $\omega_C \otimes \rho_A \succ \tau \otimes \rho'_C$, if $\rho \to \sigma$ via catalytic thermal operations. We refer to the notion $\succ$ as thermo-majorization.

Now, let us consider the scenario of preparing a pure excited state of maximum energy $\Pi^S_{\text{max}} = |E^S_{\text{max}} \rangle \langle E^S_{\text{max}}|$ from a thermal state $\tau_S$. Intuitively, if we concern ourselves only with diagonal state transformations, then this is the hardest thermal embezzling scenario possible. This is because $\Pi^S_{\text{max}} \succ \tau_S \succ \tau$.
is possible for any diagonal $\rho_S$. Therefore, whenever we investigate the case where involved states are diagonal, it suffices to analyze the preparation of such a pure excited state. The necessary and sufficient conditions are

$$\omega_C \otimes \tau_S >_T \omega_C' \otimes \Pi_{\text{max}}.$$  \hspace{1cm} (A.5)

In the next lemma, we show that given fixed Hamiltonians and dimensions, any catalyst state that succeeds in preparing such a state can also be used to facilitate any other state transformation.

**Lemma A.3 (Universal embezzlers for diagonal states).** Suppose there exists $\omega_C$, $\omega_C'$ diagonal (in $H_C$) such that $\omega_C \otimes \tau_S >_T \omega_C' \otimes \Pi_{\text{max}}$ holds, and $\| \omega_C - \omega_C' \| \leq \epsilon$. Then for any states $\rho_S$, $\rho_S'$ diagonal (in $H_S$), $\omega_C \otimes \rho_S >_T \omega_C' \otimes \rho_S'$ holds as well.

**Proof.** This can be proven by noting that $\omega_C \otimes \tau_S >_T \omega_C' \otimes \Pi_{\text{max}}$ is equivalent to the existence a thermal operation denoted by $\mathcal{M}$, such that

$$\omega_C \otimes \Pi_{\text{max}} = \mathcal{M} \omega_C' \otimes \Pi_{\text{max}}.$$  \hspace{1cm} (A.6)

is equivalent to the existence a thermal operation denoted by $\mathcal{M}$, such that

$$\omega_C \otimes \Pi_{\text{max}} = \mathcal{M} \omega_C' \otimes \Pi_{\text{max}}.$$  \hspace{1cm} (A.7)

In the next lemma, we show that given fixed Hamiltonians and dimensions, any catalyst state that succeeds in preparing such a state can also be used to facilitate any other state transformation.

**Appendix B. Optimal thermal catalyst for trivial Hamiltonians**

In this section we look at a specific thermodynamic transformation involving system ($S$) and catalyst ($C$) states of any dimension $m$ and $n = m^a$ respectively. For the trivial Hamiltonian where all states have same energy, the thermal state of the system is simply the fully mixed state $\mathbb{1}_m$, while any pure state corresponds to $\omega = \Pi_{\text{max}}$, so we simply pick $|0\rangle \langle 0|$ without loss of generality. Note that thermo-majorization conditions are reduced to the simplest form, i.e. that

$$\omega |0\rangle \langle 0| \omega' = \Pi_{\text{max}}.$$  \hspace{1cm} (B.1)

is possible if and only if the initial state majorizes the latter, i.e.

$$\omega |0\rangle \langle 0| \omega' = \Pi_{\text{max}}.$$  \hspace{1cm} (B.2)

In this section we give a construction of catalyst states which allow this transformation, and prove that our construction achieves the optimal trace distance $d(\omega_C, \omega_C') = \frac{1}{2} \| \omega_C - \omega_C' \|$ in any fixed dimension $n = m^a$. Furthermore, these states are universal embezzlers, since any catalyst which successfully creates $|0\rangle \langle 0|$ from $\omega_C$ would also allow to obtain any $\rho_S'$ from any $\rho_S$, as shown in lemma A.3.

**Definition B.1.** Consider integers $m \geq 2$ and $n = m^a$ where $a \geq 1$. Let $S_{m,n}$ be the set of $n$-dimensional catalyst state pairs $(\omega_C, \omega_C')$ enabling the transformation

$$\omega_C \otimes \frac{1}{m} \mathbb{1}_m = \omega_C' \otimes |0\rangle \langle 0|.$$  \hspace{1cm} (B.3)

Let $d_{m,n} = \min \{ d(\omega_C, \omega_C') \mid (\omega_C, \omega_C') \in S_{m,n} \}$.

**B.1. A family of catalyst states**

**Lemma B.2.** Consider a system $S$ such that $\dim(S) = m$, and a catalyst $C$ such that $\dim(C) = n = m^a$ for some integer $a \geq 1$. Consider the following catalyst state pair $(\omega_C, \omega_C')$: the state $\omega_C' = \sum_{i=1}^{n} \omega'_i |i\rangle \langle i|$, where

$$\omega'_i = \frac{1}{1 + (m - 1)a} \quad \text{and} \quad \omega_i = \omega'_i m^{a(|\omega_i| - 1)}.$$  \hspace{1cm} (B.4)
On the other hand, \( \omega_C = \sum_{i=1}^{n} \omega_i |i \rangle \langle i |, \) where
\[
\omega_i = \begin{cases} 
\omega_i^m & \text{if } i = 1, \\
\omega_i^{m-1} & \text{if } 2 \leq i \leq \frac{n}{m}, \\
0 & \text{if } i > \frac{n}{m}.
\end{cases}
\] (B.5)

Then the pair \( (\omega_C, \omega_C') \in S_{m,n} \) as defined in definition B.1, and therefore
\[
d_{m,n} \leq d(\omega_C, \omega_C') = \frac{m - 1}{1 + (m - 1)a}.
\] (B.6)

**Proof.** Before we begin the proof, it might be helpful to gain some intuition of what the structure of \( \omega_C, \omega_C' \) looks like. For \( \omega_C \), a simple way to visualize this is as follows: for the first \( m \) elements, the distribution is uniform with some probability \( \omega_1 \); for the next \( m + 1 \) up to \( m^2 \) elements the distribution is uniform again, with probability \( \omega_2 \); and so on up to \( \omega_m \). The initial \( \omega_1 \) is then chosen so that the full distribution is normalized. As for \( \omega_C' \), such a state is obtained from \( \omega_C \) by setting all the probabilities for \( i > n/m \) to be zero, while renormalizing by increasing the largest peak of the probability distribution.

Here, we prove that \( \omega_C \otimes \frac{1}{m} \omega_C' > \omega_C' \otimes [0 \rangle \langle 0 | \), the majorization relation as stated in definition B.1. It is very easy to see that this is true, once the eigenvalues of \( \omega_C \) and \( \omega_C' \) are compared
\[
\omega_C = \text{diag} \left( \begin{array}{cccccc}
\omega_1^m & \cdots & \omega_1^m & \omega_2^{m^2} & \cdots & \omega_m^{m^2} \\
1 & \cdots & 1 & 1 & \cdots & 1
\end{array} \right)
\] (B.7)
\[
\omega_C' = \text{diag} \left( \begin{array}{cccccc}
\omega_1^m & \cdots & \omega_1^m & \omega_2^{m^2} & \cdots & \omega_m^{m^2} \\
0 & \cdots & 0 & 1 & \cdots & 1
\end{array} \right).
\] (B.8)

where we have written this by making use of the fact that \( m^k = m^{k-1} = m(m^{k-1} - m^{k-2}) \).

Firstly, one can observe that since \( [0 \rangle \langle 0 | \) is a pure state with a single eigenvalue 1, \( \omega_C' \otimes [0 \rangle \langle 0 | \) has the same eigenvalues as \( \omega_C' \). On the other hand, for any two eigenvalues in \( \omega_C \), if one is greater than the other, then it is greater by at least a factor of \( m \). This implies that when we consider \( \omega_C \otimes \frac{1}{m} \), the order of these eigenvalues will not change. One can obtain the eigenvalues of \( \omega_C \otimes \frac{1}{m} \) simply by dividing each eigenvalue of \( \omega_C \) by a factor \( m \), while increasing its multiplicity also by a factor of \( m \). However, by doing so using equation (B.8), one sees that we obtain a set of eigenvalues exactly equal to those in equation (B.7). Since any vector majorizes itself, we conclude that \( \omega_C' \otimes [0 \rangle \langle 0 | \).

Note that \( \omega_1 > \omega_1' \) while \( \omega_i < \omega_i' \) for all \( i > 1 \). The trace distance between \( \omega_C \) and \( \omega_C' \) can be calculated to be
\[
d(\omega_C, \omega_C') = \frac{1}{2} \sum_{i=1}^{n} |\omega_i - \omega_i'| = \sum_{i: \omega_i > \omega_i'} \left( \omega_i - \omega_i' \right) = \omega_1 - \omega_1' = \frac{m - 1}{1 + (m - 1)a}.
\] (B.9)

This shows that
\[
d_{m,n} \leq \frac{m - 1}{1 + (m - 1)a},
\] (B.10)
since we have constructed a specific state pair achieving this trace distance. In the next section we will see that for catalysts satisfying equation (B.2), smaller values of trace distance cannot be achieved, which implies that equation (B.10) is true with equality, and the family presented above is optimal. \( \square \)

**B.2. Optimal catalysis**

In this section we show by induction that
\[
d_{m,n} \geq \frac{m - 1}{1 + (m - 1)a}.
\] (B.11)

Recall that our problem is to minimize over states \( \omega_C, \omega_C' \) the trace distance \( d(\omega_C, \omega_C') \) such that equation (B.2) is satisfied. We first show that it suffices to minimize over states which are diagonal in the same basis.
Lemma B.3 (States diagonal in the same basis). Consider fixed \( n \)-tuples of eigenvalues \((\omega_1, \ldots, \omega_n)\) and \((\omega_1', \ldots, \omega_n')\), such that \( \omega_C = \sum_i \omega_i |e_i\rangle \langle e_i| \) and \( \omega'_C = \sum_i \omega_i' |f_i\rangle \langle f_i| \) are diagonal in two different bases \( \{|e_i\}\}, \{|f_i\}\). If \((\omega_C, \omega'_C)\) satisfies equation (B.2), then there exists \( \tilde{\omega}_C = \sum_i \tilde{\omega}_i |e_i\rangle \langle e_i| \) such that \( d(\omega_C, \omega'_C) \geq d(\tilde{\omega}_C, \omega_C) \) and that \((\tilde{\omega}_C, \tilde{\omega}'_C)\) also satisfies equation (B.2).

Proof. There are two steps in this proof: firstly, we construct \( \tilde{\omega}_C \) from \( \omega'_C \) and show that the trace distance decreases by invoking data processing inequality. Then, we use Schur’s theorem to show that majorization holds. Let \( \tilde{\omega}_C = \mathcal{N}(\omega'_C) \), where \( \mathcal{N}(\rho) = \sum_i |e_i\rangle \langle e_i| \rho |e_i\rangle \langle e_i| \) is the fully dephasing channel in the basis \( \{|e_i\}\} \). Note that since \( \omega_C \) is already diagonal in \( \{|e_i\}\}, \mathcal{N}(\omega_C) = \omega_C \). Because the trace distance is non-increasing under quantum operations [26], we have

\[
d(\omega_C, \omega'_C) \geq d(\mathcal{N}(\omega_C), \mathcal{N}(\omega'_C)) = d(\omega_C, \tilde{\omega}_C).
\] (B.12)

On the other hand, we will show that \( \omega'_C \geq \tilde{\omega}_C \). For any matrix \( M \), let \( \lambda(M) \) be the vector of its eigenvalues. We want to show that \( \lambda(\omega'_C) \geq \lambda(\tilde{\omega}_C) \). Recall that \( \tilde{\omega}_C = \mathcal{N}(\omega'_C) \) and, from the definition of \( \mathcal{N} \), observe that the eigenvalues \( \lambda(\omega_C) \) are precisely the diagonal elements of \( \omega_C \) in the basis \( \{|e_i\}\} \). Schur’s theorem ([27], chapter 9, theorem B.1) says that for any Hermitian matrix \( M \), the diagonal elements of \( M \) are majorized by \( \lambda(M) \). Therefore, \( \lambda(\omega'_C) \geq \lambda(\tilde{\omega}_C) \) and thus \( \omega'_C \geq \tilde{\omega}_C \). Making use of the initial assumption \( \omega_C \otimes 1_m > \omega'_C \otimes |0\rangle \langle 0|_s \), we now see that

\[
\omega_C \otimes \frac{1}{m} 1_s > \omega'_C \otimes |0\rangle \langle 0|_s > \tilde{\omega}_C \otimes |0\rangle \langle 0|_s,
\] (B.13)

which concludes the proof. \( \square \)

We are now ready to establish our lower bound on \( d_{m,n} \) for \( n = m^a \), we will use fact established in lemma B.3, i.e. that we can take both states to be diagonal in the same basis. For the case of general \( m, n \), optimal initial/final catalyst and the corresponding trace distance can be found numerically.

Theorem B.4. Consider integers \( m \geq 2 \) and \( n = m^a \) where \( a \geq 1 \). Then

\[
d_{m,n} = \frac{m - 1}{1 + (m - 1) a},
\] (B.14)

where \( d_{m,n} \) is defined in equation (B.1). Hence, the family of catalyst states from section B.1 is optimal.

Proof. The majorization condition

\[
\omega_C \otimes \frac{1}{m} 1_s > \omega'_C \otimes |0\rangle \langle 0|_s
\] (B.15)

only depends on the eigenvalues of \( \omega \) and \( \omega' \). For fixed eigenvalues, the trace distance \( d(\omega, \omega') \) is minimized if the two states share the same eigenbasis and the eigenvalues are ordered in the same way, e.g. in decreasing order, as discussed in lemma B.3. Hence, from now on we consider only diagonal states \( \omega = \text{diag}(\omega_1, \ldots, \omega_n) \) and \( \omega' = \text{diag}(\omega_1', \ldots, \omega_n') \), where \( \omega_1 \geq \omega_2 \geq \ldots \geq \omega_n \) and \( \omega_1' \geq \omega_2' \geq \ldots \geq \omega_n' \). Here, \( \text{diag}(\cdot) \) denotes the diagonal matrix with the corresponding diagonal elements. To prove the theorem we only need to show that

\[
d_{m,n} \geq \frac{m - 1}{1 + (m - 1) a}
\] (B.16)

as the other inequality follows from the family of embezzling states exhibited in section B.1. We use induction on the power \( a \). For the base case \( a = 1 \), we need to show that \( d_{m,n} \geq 1 - 1/m \). Consider any feasible solution \((\omega, \omega')\) in dimension \( n = m \). From the majorization condition

\[
\omega \otimes \frac{1}{m} 1_m > \omega' \otimes |0\rangle \langle 0| \iff \left( \frac{\omega_1}{m}, \ldots, \frac{\omega_n}{m} \right) > \left( \frac{\omega_1'}{m}, \ldots, \frac{\omega_n'}{m}, 0, \ldots, 0 \right)
\] (B.17)

it follows that \( \omega_1/m \geq \omega_1' \) and \( \omega_i = 0 \) for \( i \geq 2 \). Hence, \( \omega_1 = 1/m \geq \omega_1' \). Since \( \omega_1' \) is the largest of the \( m \) values \( \omega_i' \), we get \( \omega_i' = 1/m \) for all \( i \). Finally, a simple calculation reveals that \( d(\omega, \omega') = 1 - 1/m \), which establishes the base case.

For the inductive step, we assume that

\[
d_{m,n} = \frac{m - 1}{1 + (m - 1) a}
\] (B.18)
for some $n = m^a$ and aim to show that
\begin{equation}
\sigma_{m,k} = \frac{m - 1}{1 + (m - 1)(a + 1)} \quad \text{B.19}
\end{equation}
for $k = m^{a+1}$. The main idea is to consider an optimal catalyst pair $(\sigma, \sigma') \in S_{m,k}$ and from it construct a catalyst pair $(\sigma, \sigma') \in S_{m,n}$ in dimension $n = m^a$. Since our construction will allow us to relate $d(\sigma, \sigma') \geq d_{m,n}$ to $d(\sigma, \sigma') = d_{m,k}$, we then obtain a lower bound on $d_{m,k}$ in terms of $d_{m,n}$, as in equation (B.18).

Let us start by using the state pair that satisfies equation (B.15) and achieves $d_{m,k}$, and from it derive some useful properties. Firstly, pick $(\sigma, \sigma') \in S_{m,k}$ so that $d(\sigma, \omega') = d_{m,k}$. As before, without loss of generality, we assume that $\sigma = \text{diag}(\omega_1, \ldots, \omega_k)$ and $\sigma' = \text{diag}(\omega'_1, \ldots, \omega'_k)$ where $\omega_1 \geq \ldots \geq \omega_k$ and $\omega'_1 \geq \ldots \geq \omega'_k$. The majorization condition
\begin{equation}
\omega \otimes \frac{1}{m} \sigma \geq \omega' \otimes |0\rangle \langle 0| \quad \text{B.20}
\end{equation}
again implies that $\omega_k \geq \omega'_k$ and $\omega_i = 0$ for $i > k/m = m^a$. To further simplify matters, we can also assume that $\omega_i \leq \omega'_i$ for all $i \geq 2$. This is because we can always replace $\omega$ with $\hat{\omega} = \text{diag}(\omega_1, \ldots, \omega_k)$, where
\begin{equation}
\hat{\omega}_i = \begin{cases} 
\omega'_i & \text{if } \omega_i > \omega'_i, \\
\omega_i & \text{otherwise},
\end{cases}
\end{equation}
for $i \geq 2$ and $\omega_i$ is chosen so that $\sum \hat{\omega}_i = 1$. In essence, all the majorization advantage of $\omega$ against $\omega'$ can be piled upon the first, largest eigenvalue of $\omega$. The reader is referred to figure B1 for a visual comparison. This replacement is valid since $(\hat{\omega}, \omega')$ still satisfies the majorization condition. Furthermore,
\begin{equation}
d(\omega, \omega') = \sum_{i: \omega_i > \omega'_i} \omega_i - \omega'_i = d(\hat{\omega}, \omega') \quad \text{B.22}
\end{equation}
implies that the distance is unchanged.

Subsequently, we proceed to bound $d_{m,n}$. To do this, construct a catalyst pair $(\sigma, \sigma') \in S_{m,n}$ in dimension $n = m^a = k/m$. Essentially, this is done by directly applying a cut to the dimension of the final catalyst state $\omega'$, reducing it to having dimension $k/m = n$. Similarly, the same amount of probability is cut from the initial state, and both states are renormalized.

Let us describe this in more detail: denote $\delta = \sum_{i \geq k/m} \omega'_i$ and pick index $s$ and value $\hat{\omega}_s \leq \omega_s$ so that $\sum_{i \geq s} \omega_i + \hat{\omega}_s = 1 - \delta$. Note that $s \leq k/m^2$, since the majorization condition equation (B.20) implies that
\begin{equation}
\sum_{i \geq k/m} \sum_{j=1}^{m} \omega_{ij} = \sum_{i \geq k/m} \omega_i \geq \sum_{i \leq k/m} \omega'_i = 1 - \delta.
\end{equation}
This inequality is obtained by summing up the first $k/m$ elements of both distributions in the lhs and rhs of equation (B.20). We now define
\begin{equation}
\sigma = \frac{1}{1 - \delta} \text{diag}(\omega_1, \ldots, \omega_{s-1}, \hat{\omega}_s, 0, \ldots, 0),
\end{equation}
\begin{equation}
\sigma' = \frac{1}{1 - \delta} \text{diag}(\omega'_1, \ldots, \omega'_{s-1}, \omega'_s, \omega'_{s+1}, \ldots, \omega'_{k/m}).
\end{equation}
Since $\sum_{i \geq s} \omega_i + \hat{\omega}_s = \sum_{i \leq k/m} \omega'_i = 1 - \delta$ the states $\sigma$ and $\sigma'$ are properly normalized. To establish that $(\sigma, \sigma') \in S_{m,n}$, we need to show that the majorization condition holds true. We consider two separate cases: when $\hat{\omega}_s = \omega_s$, and when $\hat{\omega}_s \neq \omega_s$. 

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**Figure B1.** A visual comparison between an example of states $\omega$, $\omega'$ and $\hat{\omega}$, as defined in the proof of theorem B.4. We see that whenever $\omega \geq \omega'$ (yellow bar larger than brown), we can define $\hat{\omega}$ (blue) such that $\hat{\omega} = \omega$, and renormalize $\hat{\omega}$ by increasing $\hat{\omega}_s$. Also from this visualization one can observe that the trace distance as described in equation (B.22) does not change.
If $\hat{\omega} = \omega$, then the inequalities in the majorization condition for $(\sigma, \sigma')$ have already been enforced by the majorization condition of $(\omega, \omega')$. Hence, $(\sigma, \sigma')$ is a valid catalyst pair in dimension $n = k/m$, i.e. $(\sigma, \sigma') \in S_{m,k}$. Let us now make the following two observations.

1. $d(\omega, \omega') \geq \delta$. To see this, recall that $\omega_i = 0$ for $i > k/m = n$, and thus
   \[ d(\omega, \omega') = \sum_{i=1}^{\min(k/m, n)} \omega_i - \omega_i \geq \sum_{i > k/m} \omega_i = \delta. \]  
   (B.26)

2. $d(\omega, \omega') = (1 - \delta) d(\sigma, \sigma')$. To see this, note that
   \[ \frac{d(\omega, \omega')}{1 - \delta} = \frac{1}{1 - \delta} \sum_{i=1}^{\min(k/m, n)} \omega_i - \omega_i' = \frac{\omega_i - \omega_i'}{1 - \delta} = d(\sigma, \sigma') \]  
   since only the first diagonal element of $\sigma$ is strictly larger than the corresponding diagonal element of $\sigma'$.

Combining observations 1 and 2 gives
   \[ d_{m,k} = d(\omega, \omega') = (1 - \delta) d(\sigma, \sigma') \geq \left[ 1 - d(\omega, \omega') \right] d(\sigma, \sigma') \geq (1 - d_{m,k}) d_{m,n} \]  
   (B.28)
since
   \[ d(\sigma, \sigma') \geq d_{m,n} = \frac{m - 1}{1 + (m - 1)a}. \]  
   (B.29)
Rearranging gives us
   \[ d_{m,k} \geq \frac{d_{m,n}}{1 + d_{m,n}} = \frac{m - 1}{1 + (m - 1)(a + 1)} \]  
   (B.30)
and we have completed the inductive step.

If $\hat{\omega} \neq \omega$, then the majorization inequalities involving $\hat{\omega}$ might fail to hold. Therefore, instead of $(\sigma, \sigma')$ we consider the following, slightly different, pair of states
   \[ \zeta = \sigma = \frac{1}{1 - \delta} \text{diag}(\omega_1, \ldots, \omega_{k-1}, \hat{\omega}, 0, \ldots, 0), \]  
   \[ \zeta' = \frac{1}{1 - \delta} \text{diag}(\omega_1', \ldots, \omega_{k-1}'_{m+1}, l, \ldots, \omega_n, \omega_{m+1}'_{(k-1)m}, \ldots, \omega_k'), \]  
   (B.31)  
   (B.32)
where
   \[ l = \frac{1}{m} \left( \omega_{(k-1)m+1}' + \cdots + \omega_k' \right). \]  
   (B.33)
The diagonal elements of $\zeta'$ are still in descending order, and the state is properly normalized. To argue that $(\zeta, \zeta')$ is a valid pair of catalyst states, we need to verify the majorization inequalities that are not directly implied by the majorization condition for $(\omega, \omega')$. That is, we need to verify that for all $1 \leq j \leq m$,
   \[ C + \frac{j}{m} \hat{\omega} \geq C' + jl, \]  
   (B.34)
where $C = \sum_{i=1}^{k-1} \omega_i$ and $C' = \sum_{i=1}^{(k-1)m} \omega_i'$.

We can see that this is true for the state pair $(\zeta, \zeta')$ because in this regime of equation (B.34), both sides increase linearly with the indices $j$, and for the endpoints $j = 0$ and $j = m$, the lhs is higher than the rhs, which is guaranteed by the majorization condition for $(\omega, \omega')$,
   \[ C \geq C' \quad \text{and} \quad C + \hat{\omega} \geq C' + ml. \]  
   (B.35)
Therefore, $(1 - p)C + p (C + \hat{\omega}) \geq (1 - p)C' + p (C' + ml)$ for any $0 \leq p \leq 1$. Taking $p = jm$ yields the desired inequality (B.34) and hence $(\zeta, \zeta')$ is a valid catalyst pair. Lastly, note that reasoning similar to the one in equation (B.27) can be used to deduce that
   \[ d(\omega, \omega') \frac{1}{1 - \delta} = d(\zeta, \zeta'). \]  
   (B.36)
Therefore, $d(\zeta, \zeta') = d(\sigma, \sigma')$ and we can use the argument from the previous case to complete the inductive step. By this proof of induction we have shown that $d_{m,n} \geq m - 1/(1 + (m - 1)a)$ for all $m$, $n = ma$ and $a \geq 1$. This together with the conclusion in lemma B.2 that $d_{m,n} \leq m - 1/(1 + (m - 1)a)$ proves that
\[ d_{m,n} = \frac{m - 1}{1 + (m - 1)a} , \tag{B.37} \]

and the state pair described in equation (B.4) and (B.5) is optimal. \( \square \)

### Appendix C. Limits of thermal embezzling from constraints on dimension

#### C.1. Diagonal states

In our work, we use two particular quantities, which are the Rényi divergences for \( \alpha = 1/2 \) and \( \alpha = \infty \), which for classical probability distributions have the following form:

\[ D_{1/2}(p \| q) = -2 \log \sum_i \sqrt{p_i q_i}, \quad D_{\infty}(p \| q) = \lim_{\alpha \to \infty} D_{\alpha}(p \| q) = \log \max_i \frac{p_i}{q_i} . \tag{C.1} \]

As mentioned in section A.2, given Hamiltonians \( H_S \) and \( H_C \), it suffices to consider

\[ \omega_C \otimes \tau_S \to \omega_C' \otimes \Pi_{max}^S . \tag{C.2} \]

Here, we prove whenever the dimensions of the catalyst (and system) are finite, there exists a lower bound on the accuracy of thermal embezzling. Such a bound is dependent on \( H_S \) and \( H_C \). To do so, consider the problem

\[ \min 1 \quad s. t. \quad \epsilon, \quad 0 \leq \alpha, \sigma \leq 1. \tag{C.3} \]

In [1], it was shown that for initial and target states commuting with the Hamiltonian \( H_S \), it is sufficient to consider catalyst states commuting with \( H_C \). Therefore, since \( \tau_S \) and \( \Pi_{max}^S \) both commute with \( H_C \), it is sufficient to consider initial and final catalyst states which are diagonal in the basis of \( H_C \). Since all \( \alpha \) Rényi divergences are thermal monotones according to lemma A.2, in particular the min-relative entropy \( D_{\infty} \), for \( \alpha \to \infty \),

\[ D_{\infty}(\rho \| \rho') = \max_i \log \frac{\rho_i}{\rho'_i} \tag{C.4} \]

where \( \rho_i \) and \( \rho'_i \) are the eigenvalues of \( \rho \), \( \rho' \) respectively. Therefore, satisfying the thermo-majorization conditions in equation (C.3) implies that

\[ D_{\infty}(\omega_C \otimes \tau_S \| \tau_{CS}) \geq D_{\infty}(\omega_C' \otimes \Pi_{max}^S \| \tau_{CS}) . \]

To further simplify this expression, note that \( \tau_{CS} = \tau_C \otimes \tau_S \) and that

\[ D_{\alpha}(\rho \otimes \rho' \| \sigma \otimes \sigma') = D_{\alpha}(\rho \| \sigma) + D_{\alpha}(\rho' \| \sigma'). \]

The additivity of Rényi divergences under tensor products holds for all states. Furthermore, \( D_{\alpha}(\rho \| \rho) = 0 \) for any \( \rho \). Therefore, we arrive at the expression

\[ D_{\infty}(\omega_C \| \tau_C) + 0 \geq D_{\infty}(\omega_C' \| \tau_C) + \log \frac{Z_S}{e^{-\beta E_{max}}} , \tag{C.5} \]

where \( Z_S \) is the partition function of the system. The spectral values of \( \omega_C \) and \( \omega_C' \) are denoted as \( \{ \omega_j \} \) and \( \{ \omega'_j \} \), respectively. Using the definition of \( D_{\infty} \) as shown in equation (C.1), we obtain

\[ \max_i \frac{\omega_j}{\tau_i} \geq \frac{Z_S}{e^{-\beta E_{max}}} \max_j \frac{\omega'_j}{\tau_j} \]

where

\[ \tau_j = \frac{e^{-\beta E_C}}{Z_C} \tag{C.6} \]

are the eigenvalues of the thermal state for the catalyst, for the energy eigenstate with energy eigenvalue \( E_C \), with normalization \( Z_C \), the partition function of the catalyst. Since \( \hat{\epsilon} \) is the minimum trace distance between states \( \omega_C, \omega_C' \), and \( D_{\infty} \) depends only on the maximum of \( \omega_j / \tau_j \) across the distribution, the optimal strategy to increase \( D_{\infty} \) while going from \( \omega_C' \) to \( \omega_C \) is to increase a specific \( \omega'_j \) by an amount \( \hat{\epsilon} \). Therefore, we can consider a relaxation of equation (C.3)

\[ \hat{\epsilon} = \min \quad \frac{1}{2} \quad \| \omega_C - \omega_C' \| \tag{C.7} \]

\[ \max_i \frac{\omega'_j + \hat{\epsilon}}{\tau_i} \geq \frac{Z_S}{e^{-\beta E_{max}}} \max_j \frac{\omega'_j}{\tau_j} , \quad \forall \ j, \ 0 < \omega'_j \leq 1 . \tag{C.8} \]

In the next lemma, we show that \( \epsilon \geq \hat{\epsilon} \geq \delta > 0 \) whenever \( E_{max}^C, E_{max}^S < \infty. \)
Lemma C.1 (lower bound to error in catalysis). Consider system and catalyst Hamiltonians which are finite-dimensional, and denote \( \{ E_i^S \}_{i=1}^{m_S} \), \( \{ E_i^C \}_{i=1}^{m_C} \) to be the set of energy eigenvalues respectively. Then for some fixed \( E_{\text{max}}^C \), \( E_{\text{max}}^S \), consider any probability distribution \( r \) (which corresponds to eigenvalues of a catalyst \( \omega \)), and \( \hat{e} \) such that

\[
\max_{i} \frac{r_i + \hat{e}}{\tau_i} \geq \frac{Z_S}{e^{-\beta E_{\text{max}}}} \max_{j} \frac{r_j}{\tau_j}, \quad \forall \ j, \quad 0 < r_j \leq 1,
\]

where \( \tau_i = e^{-\beta E_i^C}/Z_C \). Note that index \( i \) runs over all energy levels \( E_i^C \). Then

\[
\hat{e} \geq \left( \frac{Z_S}{e^{-\beta E_{\text{max}}} - 1} \right) e^{-\beta E_{\text{max}}} Z_C \neq 0.
\]

In other words, thermal embezzling of diagonal states with arbitrary accuracy is not possible.

**Proof.** Firstly, let \( r^*, \tau^* \) indicate the pair such that \( r^*/\tau^* = \max_{i} r_i/\tau_i \). Then

\[
\max_{i} \frac{r_i + \hat{e}}{\tau_i} \geq \max_{i} \frac{r_i + \hat{e}}{\tau_i} \geq \frac{r^*}{\tau^*} = \frac{Z_S}{e^{-\beta E_{\text{max}}} - 1}.
\]

The first term of lhs is equal to \( r^*/\tau^* \), and therefore can be grouped with the rhs to form

\[
\max_{i} \frac{\hat{e}}{\tau_i} \geq \frac{r^*}{\tau^*} \left( \frac{Z_S}{e^{-\beta E_{\text{max}}} - 1} \right) \geq \frac{Z_S}{e^{-\beta E_{\text{max}}} - 1},
\]

since we know that \( D_\infty (r \parallel q) = \log \max_{i} r_i/\tau_i = \log r^*/\tau^* \geq 0 \), therefore \( r^*/\tau^* \geq 1 \). Finally, taking the maximization of \( 1/\tau_i \) over \( i \) gives \( 1/\tau_{\text{min}} \) recall that \( \tau_i \) corresponds to probabilities of the thermal state being in the eigenstate with energy \( E_i \). Therefore, \( \tau_{\text{min}} = e^{-\beta E_{\text{min}}} / Z_C \), and we get

\[
\hat{e} \geq \left( \frac{Z_S}{e^{-\beta E_{\text{max}}} - 1} \right) e^{-\beta E_{\text{max}}} Z_C.
\]

□

C.2. Arbitrary states

The case of arbitrary states are treated separately, since our lemma A.3 on universal embezzlers hold only for diagonal states, where necessary and sufficient conditions are known for state transformations. Nevertheless, since the monotonicity of \( D_\infty \) is necessary for arbitrary state transformations \( \rho_S \to \rho'_S \), one can use techniques very similar to those in section C.1 to lower bound the embezzling error, if we minimize over diagonal catalysts.

More precisely, denote \( \hat{e} (\rho_S, \rho'_S) \) to be the solution of

\[
\min_{\omega} \frac{1}{2} \left\| \omega_C - \omega'_C \right\|_1 \quad \text{s. t.} \quad D_\infty (\omega_C \otimes \rho_S \parallel \tau_S) \geq D_\infty (\omega'_C \otimes \rho'_S \parallel \tau_S), \quad 0 \leq \omega, \sigma \leq \mathbb{1}.
\]

Recall that \( \tau_{\text{CS}} = \tau_C \otimes \tau_S \), and that \( D_\infty \) is additive under tensor products. Therefore, by defining

\[
\kappa_i (\rho_S, \rho'_S) \equiv D_\infty (\rho'_S \parallel \tau_S) - D_\infty (\rho_S \parallel \tau_S),
\]

we can rearrange the first constraint in equation (C.13)

\[
D_\infty (\omega_C \parallel \tau_C) \geq D_\infty (\omega'_C \parallel \tau_C) + \kappa (\rho_S, \rho'_S).
\]

Note that this is almost equivalent to equation (C.5), except the constant \( Z_S / e^{-\beta E_{\text{max}}} \) previously is now replaced with \( \kappa (\rho_S, \rho'_S) \). By following the same steps used to prove lemma C.1, we obtain a lower bound depending on \( \rho_S, \rho'_S \).

Lemma C.2. Consider system and catalyst Hamiltonians which are finite-dimensional, and denote \( \{ E_i^S \}_{i=1}^{m_S} \) and \( \{ E_i^C \}_{i=1}^{m_C} \) to be the set of energy eigenvalues respectively. Then for some fixed \( 0 \leq E_{\text{max}}^C \), \( E_{\text{max}}^S \), consider any probability distribution \( r \) (which corresponds to eigenvalues of a catalyst \( \omega \)), and \( \hat{e} \) such that

\[
\max_{i} \frac{r_i + \hat{e}}{\tau_i} \geq 2^\kappa (\rho_S, \rho'_S) \cdot \max_{j} \frac{r_j}{\tau_j}, \quad \forall \ j, \quad 0 < r_j \leq 1,
\]

where \( \tau_i = e^{-\beta E_i^C}/Z_C \) and \( \kappa (\rho_S, \rho'_S) = D_\infty (\rho'_S \parallel \tau_S) - D_\infty (\rho_S \parallel \tau_S) \). Note that index \( i \) runs over all energy levels \( E_i^C \). Then

\[
\hat{e} \geq \left[ 2^\kappa (\rho_S, \rho'_S) - 1 \right] e^{-\beta E_{\text{max}}}/Z_C \neq 0.
\]
This implies thermal embezzling with arbitrary accuracy, using a diagonal catalyst is not possible.

Comparing lemmas C.1 and C.2, which are very similar, one sees that for non-diagonal states lemma C.2 gives a state-dependent lower bound on the embezzling error. However for diagonal states, the bound in lemma C.1 can be made state-independent because of the existence of universal embezzlers.

C.3. Relation to energy constraints
Rather than bounding the dimension of the catalyst, one can ask if restrictions on other physical quantities such as the average energy of the catalyst would indefinitely prevent accurate embezzling from occurring. While this by itself is an independently interesting problem, we can first note that such restrictions are sometimes related to restrictions on the dimension. In one direction this is straightforward: if the catalyst is finite-dimensional, then the average energy and all other moments of energy distribution would be finite as well.

Here, we show that by restricting the first and second moments of the energy distribution of the catalyst to be finite, this implies that the states involved are always close to finite-dimensional states. In other words, if we consider the set of catalysts such that the average and variance of energy is finite, then for any such catalyst state from this set, there always exists a finite-dimensional state ε-close to it. This can be shown by invoking a simple theorem, namely the Chebyshev inequality which says that for given any finite non-zero error ε, the support of the energy distribution must be finite.

Lemma C.3 (Chebyshev inequality). Consider a random variable \( X \) with finite mean \( \bar{X} \) and finite variance \( \sigma_X^2 \), then for all \( k > 0 \),
\[
P\left[ |X - \bar{X}| \geq k \right] \leq \frac{\sigma_X^2}{k^2}.
\] (C.18)

Theorem C.4 (Chebyshev inequality applied to energy distributions). Consider a probability distribution \( p \) over some non-degenerate energy values \( E \), where both mean \( \bar{E} = \langle E \rangle \), and variance \( \sigma_E^2 = \langle (E - \bar{E})^2 \rangle \) are finite. Then for any \( \epsilon > 0 \), there exists \( E_{\text{max}} < \infty \) such that \( P\left[ E \geq E_{\text{max}} \right] \leq \epsilon \).

Proof. For any \( \epsilon > 0 \), let some \( k = \sigma_E/\sqrt{\epsilon} \). Denote \( E_{\text{max}} = \bar{E} + k \). Then by lemma C.3,
\[
P\left[ E \geq E_{\text{max}} \right] \leq P\left[ |E - \bar{E}| \geq k \right] \leq \epsilon.
\] (C.19)

Appendix D. Limits of thermal embezzling from energy constraints
In this section we provide lower bounds for the error in catalysis, given constraints on the average energy of the catalyst state. We do so by adding a constraint on the average energy of the catalyst to the problem stated in equation (C.3). By looking at the Rényi divergence for \( \alpha = 1/2 \), we can show a non-zero lower bound on the catalytic error, for cases where the partition function of the catalyst Hamiltonian \( Z_C \) is finite. This minimal assumption covers most physical scenarios, especially if we want the thermal state to be a trace class operator to begin with. Again we start with diagonal states, then later generalize to arbitrary states.

D.1. Diagonal states
Firstly, let us recall the problem stated in equation (C.3). We aim at minimizing the trace distance between all initial and final catalyst states, such that the most significant thermal embezzlement of a smaller system \( S \) can be achieved. We denote again the initial and final catalysts by \( \omega_C \) and \( \omega'_C \) with spectral values \( \{\omega_j\} \) and \( \{\omega'_j\} \).

Again, by restricting ourselves to look at the catalyst diagonal in the Hamiltonian basis, and by invoking only the thermal monotone \( D_{1/2}(\cdot \| \cdot) \), one can find the alternative relaxed problem.
\[
\begin{align*}
\min & \quad \frac{1}{2} \sum_{j=1}^{\infty} |\omega_j - \omega'_j|, \\
\text{s.t.} & \quad \sum_{j=1}^{\infty} (\omega_j^{1/2} - A^{1/2} \omega'_j^{1/2}) E_j^C \geq 0, \quad \sum_{j=1}^{\infty} \omega'_j = 1, \quad \sum_{j=1}^{\infty} \omega_j = 1, \\
& \quad \omega'_j, \omega_j \geq 0 \quad \forall \ j, \quad \text{and} \quad \sum_{j=1}^{\infty} E_j^C \omega_j \leq E, 
\end{align*}
\]

where

\[ A = \frac{Z_s}{e^{-\beta E_{\text{max}}}} \]

and \( \gamma = e^{-\beta/2} < 1 \). Furthermore, since \( A = 1/\min \tau_i \), with \( \tau_i \) forming a probability distribution (that of a thermal state), one can deduce that whenever the dimension of system \( S \) is \( m \geq 2 \), \( A \geq m \geq 2 \) holds as well.

The solution of this minimization problem serves as a lower bound to the optimal trace distance error. This problem can be relaxed to a convex optimization problem. We can arrive at a simple bound, however, with rather non-technical means. In essence, we introduce split bounds, so that the optimization can be written as two independent, individually significantly simpler optimization problems. We make use of the inequality

\[ x^{1/2} - a^{1/2} y^{1/2} \leq |x - y|^{1/2} - f(a) y, \]

which holds true for \( x, y \in [0, 1], a \geq 2 \) and with \( f: \mathbb{R}^+ \to \mathbb{R}^+ \) defined as

\[ f(a) = \frac{1}{2} \frac{a^2}{a^2 + 1}. \]

We can then relax the problem by replacing the first constraint in equation (D.1), with \( x_j \) taking the role of \( |\omega_j - \omega'_j| \), to arrive at

\[
\begin{align*}
\min & \quad \frac{1}{2} \sum_{j=1}^{\infty} x_j, \\
\text{s.t.} & \quad \sum_{j=1}^{\infty} \left[ x_j^{1/2} - f(A) \omega_j \right] e^{-\beta E_j^C/2} \geq 0, \quad \sum_{j=1}^{\infty} \omega_j = 1, \\
& \quad x_j, \omega_j \geq 0 \quad \forall \ j, \quad \text{and} \quad \sum_{j=1}^{\infty} E_j^C \omega_j \leq E. 
\end{align*}
\]

These are now two independent optimization problems, by treating \( x_j \) and \( \omega_j \) as independent variables. Define \( \epsilon_C \) to be the solution of the simple linear problem involving only variables \( \{\omega_j\} \), which we explicitly write out in corollary D.2. In this subproblem, one notes that the constraint on expectation value of the energy implies that the total probability of having relatively low energy eigenvalues cannot be vanishingly small, which we prove in lemma D.1. One can then use this fact to place a lower bound on the quantity \( \epsilon_C \), which we detail in corollary D.2.

**Lemma D.1 (lower bound to sums of eigenvalues).** Consider any probability distribution \( \{\omega_j\} \) over ascendingly ordered energy eigenvalues \( \{E_j^C\} \), with the property that the energy eigenvalues are unbounded, i.e. \( \lim_{n \to \infty} E_n^C = \infty \). If the expectation value of energy \( \sum_{j=1}^{\infty} \omega_j E_j^C \leq E \) for some finite constant \( E \), define for any \( 0 < W \leq 1 \)

\[ j(W) = \min \left\{ j : E_{j+1}^C > \frac{E}{1 - W} \right\}. \]

Then

\[ \sum_{j=1}^{j(W)} \omega_j \geq W. \]

**Proof.** One can easily prove this by contradiction. Assume that

\[ \sum_{j=1}^{j(W)} \omega_j < W \]
and therefore \( \sum_{i=j(W)+1}^{\infty} \omega_i > 1 - W \). This violates the energy constraint, since

\[
\sum_{i=j(W)+1}^{\infty} \omega_i E_i^C > (1 - W) \frac{E}{1 - W} = E.
\]

Corollary D.2 (lower bound to \( \epsilon_C \)). For a set of unbounded energy eigenvalues \( \{E_i^C\} \), consider the minimization problem

\[
\epsilon_C = \min \sum_{j=1}^{\infty} \omega_j e^{-\beta E_j^C},
\]

s. t. \( \sum_{j=1}^{\infty} \omega_j = 1 \), \( \omega_j \geq 0 \forall j \), and \( \sum_{j=1}^{\infty} E_j^C \omega_j \leq E \).

Denote \( \gamma = e^{-\beta} \in (0, 1) \). Then for \( j(W) = \min \{j: E_{j+1}^C > E/1 - W\} \),

\[
\epsilon_C \geq \max_{W \in (0,1)} W e^{j(W)},
\]

Proof. This is a direct application of lemma D.1, since the first and second constraints are satisfied automatically by any probability distribution. Given some \( W \in (0, 1) \), by lemma D.1 we know that \( \sum_{i=1}^{j(W)} \omega_i \geq W \). The objective function then can be lower bounded as

\[
\sum_{i=1}^{\infty} \omega_i e^{-\beta E_i} \geq \sum_{i=1}^{j(W)} \omega_i e^{-\beta E_i} \geq W e^{j(W)},
\]

for any such \( W \). To obtain the best lower bound, one maximizes over all \( W \in (0, 1) \).

Remark D.3 (temperature dependence). The bound obtained in corollary D.2 is dependent on temperature of the bath, and goes to zero in the limit \( T \to 0 \).

We have now solved the subproblem involving variables \( \{\omega_i\} \). Inserting the solution into the former optimisation problem, we arrive at the lower bound for \( \epsilon \),

\[
\min \frac{1}{2} \sum_{j=1}^{\infty} x_j \quad \text{s.t.} \quad \sum_{j=1}^{\infty} x_j^{1/2} e^{-\beta E_j^{1/2}} \geq f(A) \epsilon_C, \quad x_j \geq 0 \quad \forall j.
\]

The optimal solution for this minimization can easily be lower bounded by considering the Lagrange dual, which is

\[
\max - \frac{1}{4} \sum_{j=1}^{\infty} e^{-\beta E_j} \lambda^2 + \lambda f(A) \epsilon_C, \quad \text{s. t.} \lambda \geq 0.
\]

In fact, this can obviously be immediately solved as a quadratic problem in one variable. Let

\[
g(\lambda) = \sum_{j=1}^{\infty} e^{-\beta E_j} \lambda^2 + \lambda f(A) \epsilon_C,
\]

and consider the stationary point of the function by setting first derivative w.r.t. \( \lambda \) to zero,

\[
- \frac{1}{2} \sum_{i} e^{-\beta E_i} + f(A) \epsilon_C = 0,
\]

where the second derivative is negative, hence implying a maximum point. Substituting this into the objective function gives \( f(A) \epsilon_C^2/Z_C \), and hence we conclude that

\[
\epsilon \geq \frac{1}{2} \frac{f(A) \epsilon_C^2}{Z_C}.
\]

In this way, we arrive at the main result.

Theorem D.4 (energy constraint limits the accuracy of thermal catalysis). Consider the transformation\( \omega^C \otimes \tau_c \to \omega^C \otimes |E_{\text{max}}^S\rangle \langle E_{\text{max}}^S| \), where \( d_{\text{opt}} = \frac{1}{2} \| \omega^C - \omega^C \|_1 = \frac{1}{2} \epsilon \) is the error induced on the catalyst. Then for all catalyst states with finite average energy, \( d_{\text{opt}} \) is lower bounded by
\[ d_{\text{opt}} \geq \frac{1}{2} f(A)^2 \varepsilon_C^2, \]

where \( f(x) \) is defined in equation (D.4), \( A = Z \rho / e^{-\beta \varepsilon_{\text{max}}} \), \( \varepsilon_C = \max_{W \in (0,1)} W F_{E}^{C} \) and \( j(W) = \min[j : E_{j+1}^C > E/(1 - W)]. \)

In other words, thermal embezzling of diagonal states with arbitrary accuracy is not possible.

D.2. Arbitrary states

Similar to our previous discussions in section C.2, when the states \( \rho_0 \) or \( \rho'_0 \) are non-diagonal, we can still obtain a state dependent lower bound for the embezzling error. For any state \( \rho_0 \), \( \rho'_0 \), let us define the quantity

\[ \kappa_2 \left( \rho_0, \rho'_0 \right) := D_{1/2} \left( \rho_0 ^{\otimes} \| \tau_2 \right) - D_{1/2} \left( \rho'_0 ^{\otimes} \| \tau_2 \right), \]  

(D.14)

Then a lower bound can be obtained by following the steps as proved in section D.1, only now replacing the constant \( A \) defined in equation (D.2) with a state-dependent function.

**Lemma D.5.** For arbitrary states \( \rho_0 \) and \( \rho'_0 \), consider the transformation \( \omega_0 \otimes \rho_0 \rightarrow \omega'_0 \otimes \rho'_0 \), where

\[ d_{\text{opt}} = \frac{1}{2} \| \omega_0 - \omega'_0 \| = \frac{1}{2} e \]

is the error induced on the catalyst. Then for all catalyst states with finite average energy, \( d_{\text{opt}} \) is lower bounded by

\[ d_{\text{opt}} \geq \frac{1}{2} f \left( 2\varepsilon_2 (\rho_0, \rho'_0) \right)^2 \varepsilon_C^2, \]

where \( f(x) \) is defined in equation (D.4), \( \kappa_2 \left( \rho_0, \rho'_0 \right) = D_{1/2} \left( \rho_0 ^{\otimes} \| \tau_2 \right) - D_{1/2} \left( \rho'_0 ^{\otimes} \| \tau_2 \right), \varepsilon_C = \max_{W \in (0,1)} W F_{E}^{C} \) and \( j(W) = \min[j : E_{j+1}^C > E/(1 - W)]. \) This implies that thermal embezzling with arbitrary accuracy using a diagonal catalyst is not possible.

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