Location-Domination and Matching in Cubic Graphs

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Abstract

A dominating set of a graph $G$ is a set $D$ of vertices of $G$ such that every vertex outside $D$ is adjacent to a vertex in $D$. A locating-dominating set of $G$ is a dominating set $D$ of $G$ with the additional property that every two distinct vertices outside $D$ have distinct neighbors in $D$; that is, for distinct vertices $u$ and $v$ outside $D$, $N(u) \cap D \neq N(v) \cap D$ where $N(u)$ denotes the open neighborhood of $u$. A graph is twin-free if every two distinct vertices have distinct open and closed neighborhoods. The location-domination number of $G$, denoted $\gamma_L(G)$, is the minimum cardinality of a locating-dominating set in $G$. It is conjectured [D. Garijo, A. González and A. Márquez. Resolving sets for breaking symmetries of graphs. arXiv pre-print:1401.3686, 2014] that if $G$ is a twin-free graph of order $n$ without isolated vertices, then $\gamma_L(G) \leq \frac{n}{2}$. We prove the conjecture for cubic graphs. We rely heavily on proof techniques from matching theory to prove our result.

Keywords: Locating-dominating set; Dominating set; Matching.
AMS subject classification: 05C69

1 Introduction

A dominating set in a graph $G$ is a set $D$ of vertices of $G$ such that every vertex outside $D$ is adjacent to a vertex in $D$. The domination number, $\gamma(G)$, of $G$ is the minimum cardinality of a dominating set in $G$. The literature on the subject of domination parameters in graphs up to the year 1997 has been surveyed and detailed in the two books [10, 11]. In this paper, we focus our attention on a variation of domination, called location-domination, which is widely studied in the literature. A locating-dominating set is a dominating set $D$ that locates all the vertices in the sense that every vertex outside $D$ is uniquely determined by its neighborhood in $D$. The location-domination number of $G$, denoted $\gamma_L(G)$, is the minimum cardinality of a locating-dominating set in $G$. The

*Research supported in part by the South African National Research Foundation and the University of Johannesburg.
concept of a locating-dominating set was introduced and first studied by Slater [16,17] and studied in [8,11,15,16,17,18] and elsewhere.

A classic result due to Ore [14] states that every graph without isolated vertices has a dominating set of cardinality at most one-half its order. As observed in [8], while there are many graphs (without isolated vertices) which have location-dominination number much larger than one-half their order, the only such graphs that are known contain many twins, that is, pairs of vertices with the same closed or open neighborhood. Garijo et al. [9] therefore conjectured that in the absence of twins, the classic bound of one-half the order for the domination number also holds for the location-dominination number.

**Conjecture 1** (Garijo et al. [9]). Every twin-free graph $G$ of order $n$ without isolated vertices satisfies $\gamma_L(G) \leq \frac{n}{2}$.

The same authors prove Conjecture 1 for graphs without 4-cycles (which include trees) and for the class of graphs with independence number at least one-half the order (which includes bipartite graphs). In [8], the authors provide several constructions for twin-free graphs with location-dominination number one-half their order. The variety of these constructions shows that these graphs have a rich structure, which is an indication that the conjecture might be difficult to prove. Further support is given to the conjecture in [8] where it is proved for split graphs and co-bipartite graphs, and in [7] where it is proved for line graphs. The following theorem summarizes the known results about Conjecture 1.

**Theorem** ([8,9,12]) Conjecture 1 is true if the twin-free graph $G$ of order $n$ (without isolated vertices) satisfies any of the following conditions.

(a) ([9]) $G$ has no 4-cycles.
(b) ([9]) $G$ has independence number at least $\frac{n}{2}$.
(c) ([9]) $G$ has clique number at least $\lceil \frac{n}{2} \rceil + 1$.
(d) (8) $G$ is a split graph or a co-bipartite graph.
(e) (7) $G$ is a line graph.
(f) (12) $G$ is a cubic claw-free graph.

In this paper, we continue to advance the study of the conjecture by proving it for the class of cubic graphs, as stated in our main theorem:

**Theorem 2.** If $G$ is a cubic twin-free graph of order $n$, then $\gamma_L(G) \leq \frac{n}{2}$.

The proof of Theorem 2 extensively uses matchings in graphs. We start by giving graph theory definitions and notations in Section 2 and we prove Theorem 2 in Section 3.

## 2 Definitions and notation

For notation and graph theory terminology, we in general follow [11]. Specifically, let $G$ be a graph with vertex set $V(G)$, edge set $E(G)$ and with no isolated vertex. The open neighborhood of a vertex $v \in V(G)$ is $N_G(v) = \{u \in V | uv \in E(G)\}$ and its closed neighborhood is the set $N_G[v] = N_G(v) \cup \{v\}$. The degree of $v$ is $d_G(v) = |N_G(v)|$. If the graph $G$ is clear from the context, we simply write $V$, $E$, $N(v)$, $N[v]$ and $d(v)$ rather than $V(G)$, $E(G)$, $N_G(v)$, $N_G[v]$ and $d_G(v)$, respectively. Given a set $S$ of edges, we will denote by $G - S$ the subgraph obtained from $G$ by
deleting all edges of \( S \). For a set \( S \) of vertices, \( G - S \) is the graph obtained from \( G \) by removing all vertices of \( S \) and removing all edges incident to vertices of \( S \). The subgraph induced by \( S \) is denoted by \( G[S] \). A cycle on \( n \) vertices is denoted by \( C_n \) and a path on \( n \) vertices by \( P_n \). An odd component of \( G \) is a component of \( G \) of odd order. The number of odd components of \( G \) is denoted by \( \text{oc}(G) \).

A set \( D \) is a dominating set of \( G \) if \( N[v] \cap D \neq \emptyset \) for every vertex \( v \) in \( G \), or, equivalently, \( N[S] = V(G) \). Two distinct vertices \( u \) and \( v \) in \( V(G) \setminus D \) are located by \( D \) if they have distinct neighbors in \( D \); that is, \( N(u) \cap D \neq N(v) \cap D \). If a vertex \( u \in V(G) \setminus D \) is located from every other vertex in \( V(G) \setminus D \), we simply say that \( u \) is located by \( D \). For \( k \geq 1 \) if \( X \) is a set of vertices in \( G \) and \( x \in V(G) \setminus X \), then the vertex \( x \) is said to be \( k \)-dominated by \( X \) if \( x \) has exactly \( k \) neighbors inside \( X \); that is, \( |N(x) \cap X| = k \).

A set \( S \) is a locating set of \( G \) if every two distinct vertices outside \( S \) are located by \( S \). In particular, if \( S \) is both a dominating set and a locating set, then \( S \) is a locating-dominating set. Further, if \( S \) is both a total dominating set and a locating set, then \( S \) is a locating-total dominating set (where \( S \) is a total dominating set of \( G \) if every vertex of \( G \) is adjacent to some vertex in \( S \)). We remark that the only difference between a locating set and a locating-dominating set in \( G \) is that a locating set might have a unique non-dominated vertex.

An independent set in \( G \) is a set of vertices no two of which are adjacent. Two distinct edges in a graph \( G \) are independent if they are not adjacent in \( G \) (i.e., the two edges are not incident with a common vertex). A set of pairwise independent edges of \( G \) is called a matching in \( G \). A matching of maximum cardinality in \( G \) is called a maximum matching in \( G \). The number of edges in a maximum matching of a graph \( G \) is called the matching number of \( G \), denoted by \( \alpha'(G) \). Let \( M \) be a specified matching in a graph \( G \). A vertex \( v \) of \( G \) is an \( M \)-matched vertex if \( v \) is incident with an edge of \( M \); otherwise, \( v \) is an \( M \)-unmatched vertex. If the matching \( M \) is clear from context, we simply call a \( M \)-matched vertex a matched vertex and a \( M \)-unmatched vertex an unmatched vertex.

Two distinct vertices \( u \) and \( v \) of a graph \( G \) are open twins if \( N(u) = N(v) \) and closed twins if \( N[u] = N[v] \). Further, \( u \) and \( v \) are twins in \( G \) if they are open twins or closed twins in \( G \). A graph is twin-free if it has no twins. We use the standard notation \([k] = \{1, 2, \ldots, k\}\).

## 3 Proof of Theorem 2

In this section, we present a proof of Theorem 2. Our proof relies heavily on matching theory in graphs. We begin with some useful definitions and lemmas related to matchings.

### 3.1 Useful definitions and lemmas

We shall need the following theorem of Berge [1] about the matching number of a graph, which is sometimes referred to as the Tutte-Berge formulation for the matching number. Recall that \( \text{oc}(G) \) denotes the number of odd components in a graph \( G \).

**Theorem 3.** (Tutte-Berge Formula) For every graph \( G \),

\[
\alpha'(G) = \min_{X \subseteq V(G)} \frac{1}{2} \frac{1}{2} (|V(G)| + |X| - \text{oc}(G - X)).
\]
We shall also need the following structural result about maximum matchings in graphs which is a consequence of the proof of the Tutte-Berge Formula.

**Theorem 4.** (1) Let \( G = (V, E) \) be a graph and let \( X \) be a proper subset of vertices of \( G \) such that \((|V| + |X| - oc(G - X))/2 \) is minimum. If \( M \) is a maximum matching in \( G \), then \(|M| = (|V| + |X| - oc(G - X))/2 \) and there are exactly \( oc(G - X) - |X| \) vertices that are \( M \)-unmatched. Furthermore, if \( M_X \) is the subset of edges of \( M \) that belong to \( G - X \), then every vertex in \( G - X \) is \( M_X \)-matched, except for exactly one vertex from each odd component of \( G - X \). If \( U \) denotes this set of \( oc(G - X) \) vertices that are \( M_X \)-unmatched, one from each odd component of \( G - X \), then \( X \) is \( M \)-matched to a subset of vertices in \( U \).

The structure described in Theorem 4 is illustrated in Figure 1.

**Definition of the set \( D_G(M) \).** Let \( G \) be a graph and let \( M \) be a maximum matching of \( G \). We define \( D_G(M) \) to be the collection of all sets \( D \) such that the following holds:

- For every edge \( uv \in M \), if exactly one of \( u \) and \( v \) has a neighbor that is \( M \)-unmatched, then the vertex in \( \{u, v\} \) with an \( M \)-unmatched neighbor belongs to \( D \).
- For every edge \( uv \in M \), if neither \( u \) nor \( v \) has an \( M \)-unmatched neighbor or if both \( u \) and \( v \) have a (common) \( M \)-unmatched neighbor, then exactly one of \( u \) and \( v \) belongs to \( D \).

If the graph \( G \) is clear from the context, we simply write \( D(M) \) rather than \( D_G(M) \).

**Definition of a \( D \)-bad pair.** Given a set \( D \subseteq V(G) \), we define a \( D \)-bad pair of vertices as two vertices in \( V(G) \setminus D \) that are not located by \( D \). If the set \( D \) is clear from the context, we simply write that a pair of vertices is a bad pair rather than a \( D \)-bad pair.
In our proof, we will use the following lemmas.

**Lemma 5.** Let $G$ be a cubic graph, let $M$ be a maximum matching of $G$, and let $D \in \mathcal{D}_G(M)$. Then, $D$ is a dominating set of $G$, and each $M$-unmatched vertex is dominated by at least two vertices of $D$.

**Proof.** It follows readily from the two properties of sets $D \in \mathcal{D}_G(M)$, that every $M$-matched vertex is dominated by $D$. If $x$ is an $M$-unmatched vertex, then since $G$ is cubic, the vertex $x$ is adjacent to two $M$-matched vertices that are incident with distinct edges, $e_1$ and $e_2$, say, of $M$. Hence, by the construction of $D$, the set $D$ contains a neighbor of $x$ incident with $e_1$ and a neighbor of $x$ incident with $e_2$. Thus, $x$ is dominated by at least two vertices of $D$.

**Lemma 6.** Let $G$ be a twin-free cubic graph, let $M$ be a maximum matching of $G$, and let $D \in \mathcal{D}_G(M)$. Then, the vertices of each $D$-bad pair are $2$-dominated.

**Proof.** Let $\{u, v\}$ be a $D$-bad pair. If $u$ and $v$ were $3$-dominated by $D$, then they would be open twins, a contradiction. Hence, $u$ and $v$ are dominated by at most two vertices of $D$. By Lemma 5, only $M$-matched vertices can be $1$-dominated by $D$, but if $u$ and $v$ are $M$-matched vertices, they would each be dominated by the vertex of $D$ that they are matched to under $M$ and would therefore not form a $D$-bad pair, a contradiction. Hence, $u$ and $v$ are $2$-dominated by $D$.

**Lemma 7.** Let $G$ be a twin-free cubic graph. Among all maximum matchings $M$ of $G$ and all sets $D \in \mathcal{D}_G(M)$, let the matching $M_0$ and the set $D_0 \in \mathcal{D}_G(M_0)$ be chosen so that the number of $D_0$-bad pairs is a minimum. Then, the vertices of each $D_0$-bad pair are $M_0$-matched vertices.

**Proof.** Let $X$ be a proper subset of vertices of $G$ such that $(|V| + |X| - \text{oc}(G - X))/2$ is minimum. The structure of the graph $G$ with respect to the matching $M_0$ and the set $X$ is described in Theorem 4. Let $\{u, v\}$ be an arbitrary $D_0$-bad pair. Suppose to the contrary that they are not both $M_0$-matched vertices. By Lemma 6, both $u$ and $v$ are $2$-dominated by $D_0$. By the definition of a set in $\mathcal{D}_G(M_0)$, if both $u$ and $v$ are $M_0$-unmatched and $2$-dominated, they would be open twins, a contradiction. Therefore, exactly one of $u$ and $v$ is $M_0$-unmatched. Renaming $u$ and $v$ if necessary, we may assume that $u$ is $M_0$-unmatched. Thus, by Theorem 4, the vertex $u$ belongs to an odd component, $C_u$, say, of $G - X$. Let $x$ and $y$ be the two common neighbors of $u$ and $v$ in $D_0$. Let $x'$ and $y'$ be the vertices $M_0$-matched to $x$ and $y$, respectively.

Suppose that both $x$ and $y$ belong to $C_u$. If $v \notin V(C_u)$, then $v \in X$. But then the vertex that is $M_0$-matched to $v$ belongs to $D_0$, implying that $v$ would be $3$-dominated by $D_0$, contradicting Lemma 5. Hence, $v \in V(C_u)$. If $v$ is matched to neither $x$ nor $y$ by $M_0$, then, once again, $v$ would be $3$-dominated by $D_0$, a contradiction. Hence, $v$ is $M_0$-matched to either $x$ or $y$. Renaming $x$ and $y$ if necessary, we may assume that $xv \in M_0$, and so $v = x'$. If $u$ and $v$ are adjacent, then $u$ and $v$ would be closed twins, a contradiction. Hence, $u$ and $v$ are not adjacent. The third neighbor of $u$, different from $x$ and $y$, is therefore the vertex $y'$ that is $M_0$-matched to $y$ (for otherwise, by the definition of $\mathcal{D}_G(M)$, $u$ would be $3$-dominated, a contradiction). We now consider the set $D' = (D_0 \setminus \{y\}) \cup \{y'\}$. We note that $y$ and $y'$ have a common $M_0$-unmatched neighbor, namely $u$, implying that $D' \in \mathcal{D}(M_0)$. Suppose there is a vertex $z$ different from $u$ that is adjacent to both $x$ and $y'$. Then, $z$ is either in $X$ or in $C_u$. In both cases, $z$ is $M_0$-matched. If $z = v$, then $u$ and $z$ are open twins, a contradiction. Since $N(y) = \{u, v, y'\}$ while $z$ is adjacent to $x$, we note that $z \neq y$. Therefore, $z \notin \{v, y\}$ and the $M_0$-matched neighbor of $z$ is in $D'$, implying that $z$ is $3$-dominated by $D'$. Hence, the vertex $u$ is the only vertex dominated only by $x$ and $y'$ in $D'$ and it is therefore located by $D'$. Moreover, both $v$ and $y$ are $1$-dominated by $D'$, and are therefore located by $D'$ by Lemma 6. Finally, no other vertex has been affected by the removal of $y$ from $D_0$. Hence, the
number of $D'$-bad pairs is strictly less than the number of $D_0$-bad pairs, contradicting our choice of $D_0$. Therefore, at most one of $x$ and $y$ belong to $C_u$.

Suppose that exactly one of $x$ and $y$ belong to $C_u$. Renaming $x$ and $y$ if necessary, we may assume that $x \in V(C_u)$. Then, $y \in X$. If $v \in X$ or if $v \in V(C_u) \setminus \{x\}$, then $v$ would be $3$-dominated by $D_0$, a contradiction. Hence, $v = x'$, and so $v$ is $M_0$-matched to $x$. Since $u$ is $2$-dominated by $D_0$, the vertex $u$ is adjacent to either $v$ or $y'$. Since $y'$ belongs to a component of $G - X$ different from $C_u$, the vertex $u$ is adjacent to $v$. But then $u$ and $v$ are closed twins, a contradiction.

Therefore, both $x$ and $y$ belong to $X$. This implies that $u$, $x'$ and $y'$ belong to three different components of $G - X$. In particular, $u$ is adjacent to neither $x'$ nor $y'$, implying that the third neighbor of $u$ different from $x$ and $y$ is an $M_0$-matched vertex and therefore belongs to the set $D_0$ by the construction of sets in $D_G(M_0)$. Thus, $u$ is then $3$-dominated by $D_0$, a contradiction. This completes the proof of the lemma. □

Note that in any cubic twin-free graph $G$, every 4-cycle is an induced 4-cycle. The following structure will play an important role in our proof.

**Definition of a bad $(D, M)$-matched 4-cycle.** Let $C: u_C u'_C v_C v'_C u_C$ be a 4-cycle in a (cubic twin-free) graph $G$, $M$ a matching of $G$, and $D$, a subset of vertices of $G$. We say that $C$ is a bad $(D, M)$-matched 4-cycle if $u_C u'_C \in M$, $v_C v'_C \in M$, $D \cap V(G) = \{u'_C, v'_C\}$ and $v_C$ is adjacent to exactly two vertices of $D$ (and so, $N(v_C) \cap D = \{u'_C, v'_C\}$).

Given two bad $(D, M)$-matched 4-cycles, $A$ and $B$, we say that $A$ is dependent on $B$ via the vertex $u'_A$ or $v'_A$ if $u_B$ is adjacent to $u'_A$ or to $v'_A$, respectively. An illustration is given in Figure 2. We note that if $A$ is dependent on $B$, then $u_B$ is 3-dominated by $D$.

![Figure 2: Two bad $(D, M)$-matched 4-cycles $A$ and $B$, where $A$ is dependent on $B$ via $u'_A$. Edges of $M$ are thickened. Black vertices belong to $D$, and white vertices do not.](image)

Given a set $S$ of vertex-disjoint bad $(D, M)$-matched 4-cycles of a graph $G$, let $\overline{G}(S)$ be the digraph with vertex set $S$ and where $(A, B)$ is an arc in $\overline{G}(S)$ if $A$ is dependent on $B$. We remark that since $G$ is cubic, every vertex in $\overline{G}(S)$ has out-degree at most 2. Further by definition of a bad $(D, M)$-matched 4-cycle, every vertex in $\overline{G}(S)$ has in-degree at most 1.

Given a rooted tree $T$ with root $r$, by an orientation of $T$ we mean orienting every arc of $T$ from a parent to its child.

### 3.2 Proof of Main Result

We are now in a position to prove our main result, namely Theorem[2]. Recall its statement.

**Theorem[2]** If $G$ is a cubic twin-free graph of order $n$, then $\gamma_L(G) \leq \frac{n}{2}$. 

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Proof of Theorem 2. Among all maximum matchings $M$ of $G$ and all sets $D \in \mathcal{D}_G(M)$, we choose the matching $M_0$ and the set $D_0 \in \mathcal{D}_G(M_0)$ so that the number of $D_0$-bad pairs is a minimum. Let $X$ be a proper subset of vertices of $G$ such that $(|V| + |X| - \text{occ}(G - X))/2$ is minimum. The structure of the graph $G$ with respect to the matching $M_0$ and the set $X$ is described in Theorem 2.

We now describe the structure of $D_0$-bad pairs:

Claim A. Every $D_0$-bad pair $\{u, v\}$ belongs to a common bad $(D_0, M_0)$-matched 4-cycle, say $R$. Further, there exists a set $S_{u,v}$ of vertex-disjoint bad $(D_0, M_0)$-matched 4-cycles containing $R$ such that the following holds:

(a) For every 4-cycle $C \in S_{u,v}$ and every vertex $x \in \{u'_C, v'_C\}$, either $x$ is adjacent to an $M_0$-unmatched vertex in $G$, or $C$ is dependent on some other $C' \in S_{u,v}$ via the vertex $x$.

(b) $G(S_{u,v})$ is an oriented tree rooted at $R$.

(c) For every 4-cycle $C \in S_{u,v}$, if both $u'_C$ and $v'_C$ have an $M_0$-unmatched neighbor, then these neighbors are distinct and $\{u'_C, v'_C\} \subseteq X$.

Proof of Claim A. Let $\{u, v\}$ be a $D_0$-bad pair. Thus, $u$ and $v$ are vertices outside $D_0$ that are not located by $D_0$. By Lemma 6, both $u$ and $v$ are 2-dominated by $D_0$, and by Lemma A, both $u$ and $v$ are $M_0$-matched. Let $u'$ and $v'$ be the $M_0$-matched neighbors of $u$ and $v$, respectively. Since $u$ and $v$ are 2-dominated by $D_0$, we note that $u'$ and $v'$ are the two common neighbors of $u$ and $v$ in $D_0$. Thus, $C_{uv}: uu'vv'u$ is a bad $(D_0, M_0)$-matched 4-cycle in $G$. As observed earlier, every 4-cycle in $G$ is an induced 4-cycle. Hence, let $x, y, u''$ and $v''$ be the neighbors of $u, v, u'$ and $v'$, respectively, that do not belong to this 4-cycle $C_{uv}$. Let $D_u = (D_0 \setminus \{u\}) \cup \{v\}$ and let $D_v = (D_0 \setminus \{v\}) \cup \{u\}$. We proceed further with the following series of subclaims.

Claim A.1. The following holds.

(a) If $D_u \notin \mathcal{D}_G(M_0)$, then $u''$ is an $M_0$-unmatched vertex that is not adjacent to $u$.

(b) If $D_u \in \mathcal{D}_G(M_0)$, then the only $D_u$-bad pair that is not a $D_0$-bad pair is $\{u'', z\}$ for some vertex $z$. Moreover, $u''$ and $z$ are part of a bad $(D_0, M_0)$-matched 4-cycle $C$ of $G$, and $C_{uv}$ is dependent on $C$ via the vertex $u'$.

Proof of Claim A.1. By definition of the sets in the family $\mathcal{D}_G(M_0)$, if $D_u \notin \mathcal{D}_G(M_0)$, then $u''$ is an $M_0$-unmatched vertex and $u''$ is not adjacent to $u$, proving Part (a) of the statement.

To prove Part (b), suppose that $D_u \in \mathcal{D}_G(M_0)$. By our choice of $M_0$ and $D_0$, there are at least as many $D_u$-bad pairs as $D_0$-bad pairs. The only vertices that could potentially be negatively affected (in the sense that they are located by $D_0$ but not by $D_u$) by removing $u'$ from $D_0$ and replacing it with the vertex $u$ are $u'$, $u''$ and $v$. By Lemma 5, two vertices forming a $D_{u'}$-bad pair are 2-dominated by $D_u$. The vertex $v$ is 1-dominated by $D_{u'}$, and hence it is located by $D_u$.

Suppose that $u'$ is not located by $D_u$ from some other vertex outside $D_u$. Then, this vertex must be $x$, the neighbor of $u$ not on $C_{uv}$. Considering the $D_{u'}$-bad pair $\{u', x\}$, and noting that $u'$ and $x$ are 2-dominated by $D_u$, we deduce that $u'' \in D_u$. If $x$ is $M_0$-unmatched, then by the definition of $\mathcal{D}_G(M_0)$, we would have $u \in D_0$ and $u' \notin D_0$, a contradiction. Hence, $x$ is $M_0$-matched and its matched neighbor is in $D_u$. Since $x$ is 2-dominated, we have $xu'' \in M_0$. We now consider the maximum matching $M' = (M_0 \setminus \{u', v''\}) \cup \{u''', vv''\}$, and we let $D' = (D_0 \setminus \{u'\}) \cup \{v\}$. We note that $u''x \in M'$. Since neither $u$ nor $u'$ has an $M'$-unmatched neighbor, we note that $D' \in \mathcal{D}(M')$. The only vertices that could potentially be negatively affected (in the sense that they are located by $D_0$ but not by $D'$) by these changes are the vertices dominated by $u'$ in $D_0$ and that do not belong to $D'$. The only such vertices are $u$ and $u'$. By Lemma 5, if two vertices form a $D'$-bad pair, then they are 2-dominated by $D'$. The vertex $u$ is 1-dominated by $D'$, and hence it is located by $D'$.
Thus, the vertex \( u' \) is not located by \( D' \) from some other vertex outside \( D' \). Such a vertex must be adjacent to both \( v \) and \( u'' \), and is therefore the neighbor of \( v \) outside \( C_{uv} \), namely the vertex \( y \). If \( x = y \), then \( u \) and \( v \) would be open twins, a contradiction. Hence, \( x \neq y \). If \( y \) is \( M' \)-matched, since \( u''x \in M' \) and \( u''v \in M' \), the vertex \( y \) is 3-dominated by \( D' \), a contradiction. Hence, \( y \) is \( M' \)-unmatched (and \( M \)-unmatched). Then, since \( D_0 \in D_G(M_0) \), the vertices \( y \) and \( v' \) are adjacent. Hence, \( y \) is 3-dominated by \( D' \), a contradiction. Thus, \( u' \) is located by \( D_u \).

Therefore, among the vertices dominated by \( u' \) in \( D_0 \), the vertex \( u'' \) is the only vertex that was located by \( D_0 \) but that is not located by \( D_u \) from some other vertex, \( z \) say, outside \( D_u \). Thus, \( \{u'', z\} \) is the only pair of vertices located by \( D_0 \) but not by \( D_u \). Hence, the number of \( D_u \)-bad pairs is the same as the number of \( D_0 \)-bad pairs (since \( \{u, v\} \) is not a \( D_u \)-bad pair) and we can apply Lemma 6 to \( D_u \) to deduce that \( u'' \) and \( z \) are \( M_0 \)-matched vertices. By Lemma 6, both \( u'' \) and \( z \) are 2-dominated by \( D_u \). Let \( w \) and \( t \) be the \( M_0 \)-matched neighbors of \( u'' \) and \( z \), respectively. Since \( u'' \) and \( z \) are 2-dominated by \( D_u \) and \( D_u \in D_G(M_0) \), we note that \( w \) and \( t \) are the two common neighbors of \( u'' \) and \( z \) in \( D_u \). Thus, the 4-cycle \( C: u''wztu'' \) is a bad \( (D_0, M_0) \)-matched 4-cycle in \( G \) with \( u'' = u_C, w = u'_C, z = v_C \) and \( t = v'_C \), and \( C_{uv} \) is dependent on \( C \) via the vertex \( u' \). This establishes Part (b) of the statement. (c)

Interchanging the roles of \( u \) and \( v \) in the proof of Claim A.1, we have the following analogous result for the vertex \( v \).

Claim A.2. The following holds.
(a) If \( D_v \notin D_G(M_0) \), then \( v' \) is an \( M_0 \)-unmatched vertex that is not adjacent to \( v \).
(b) If \( D_v \in D_G(M_0) \), then the only \( D_v \)-bad pair that is not a \( D_0 \)-bad pair is \( \{v'', z\} \) for some vertex \( z \). Moreover, \( v'' \) and \( z \) are part of a bad \( (D_0, M_0) \)-matched 4-cycle \( C \) of \( G \), and \( C_{uv} \) is dependent on \( C \) via the vertex \( v' \).

Let \( R \) denote the bad \( (D_0, M_0) \)-matched 4-cycle \( C_{uv}: u'u'v'u \), where \( u = u_R, u' = u'_R, v = v_R \) and \( v' = v'_R \). We now show the existence of a set \( S_{u,v} \) of vertex-disjoint bad \( (D_0, M_0) \)-matched 4-cycles containing \( R \) such that conditions (a) and (b) in the statement of Claim A hold. If both \( u' \) and \( v' \) have an \( M_0 \)-unmatched neighbor, then we let \( S_{u,v} = \{R\} \) and we are done.

Otherwise, renaming vertices if necessary, we may assume, by Claim A.1 and Claim A.2, that \( D_u \in D_G(M_0) \) and that \( R \) is dependent on a bad \( (D_0, M_0) \)-matched 4-cycle \( C \) via vertex \( u'_R \). Now, since by Claim A.1(b) the number of \( D_u \)-bad pairs is the same as the number of \( D_0 \)-bad pairs (hence \( D_u \) also minimizes the number of bad pairs), we can apply Claim A.1 and Claim A.2 to \( C, D_u \) and to the \( D_u \)-bad pair \( \{u_C, v_C\} \). This shows that each of \( u'_C \) and \( v'_C \) either have an \( M_0 \)-unmatched neighbor, or \( C \) is dependent on some other bad \( (D_0, M_0) \)-matched 4-cycle via this vertex. Repeating this process as long as possible yields a set \( S_{u,v} \) of bad \( (D_0, M_0) \)-matched 4-cycles, where for each bad \( (D_0, M_0) \)-matched 4-cycle in \( S_{u,v} \), different from \( R \) there is some other bad \( (D_0, M_0) \)-matched 4-cycle in \( S_{u,v} \) that depends on it, and satisfies the properties in Claim A.1 and Claim A.2. This establishes Property (a) in the statement of Claim A. Moreover we have the following.

Claim A.3. Any two distinct 4-cycles in \( S_{u,v} \) are vertex-disjoint.

Proof of Claim A.3. Let \( A : u_A u'_A v_A v'_A \) and \( B : u_B u'_B v_B v'_B \) be two distinct 4-cycles of \( S_{u,v} \). If they have a common vertex, since their vertices are pairwise \( M_0 \)-matched, they must have two vertices in common, and these vertices must be \( M_0 \)-matched to each other. But then, the vertex that belongs to both \( A \) and \( B \) but does not belong to \( D_0 \) must be 3-dominated by \( D_0 \), a contradiction. (c)

Now, consider the digraph \( G(S_{u,v}) \), which by Claim A.3 is well-defined. The following properties
hold in the digraph $\overrightarrow{G}(S_{u,v})$. Recall that the distance from a vertex $x$ to a vertex $y$ in a directed graph $D$ is the minimum length among all directed paths from $x$ to $y$ in $D$.

**Claim A.4.** The following holds.

(a) $R$ has in-degree 0 in $\overrightarrow{G}(S_{u,v})$.

(b) Every vertex in $\overrightarrow{G}(S_{u,v})$ different from $R$ has in-degree exactly 1.

(c) $\overrightarrow{G}(S_{u,v})$ has no directed cycle.

**Proof of Claim A.4.** To see that Part (a) holds, observe that if some bad $(D_0, M_0)$-matched 4-cycle $C$ was dependent on $R$ say, on vertex $u_R$, then $u_R$ would be 3-dominated by $D_0$, a contradiction.

For Part (b), we show firstly that $\overrightarrow{G}(S_{u,v})$ has maximum in-degree 1. Suppose to the contrary that for some bad $(D_0, M_0)$-matched 4-cycle $A$ in $S_{u,v}$, there are two other bad $(D_0, M_0)$-matched 4-cycles $B$ and $C$ of $S_{u,v}$ that are both dependent on $A$. Since $B$ is dependent on $A$, the vertex $u_A$ is adjacent to $u_B$ or $v_B$. Since $C$ is dependent on $A$, the vertex $u_A$ is adjacent to $u_C$ or $v_C$. Thus, the vertex $u_A$ has degree at least 4, a contradiction. We show secondly that $R$ is the only vertex in $\overrightarrow{G}(S_{u,v})$ with in-degree 0. As observed in the paragraph immediately preceding Claim A.3, for each bad $(D_0, M_0)$-matched 4-cycle in $S_{u,v}$ different from $R$ there is some other bad $(D_0, M_0)$-matched 4-cycle in $S_{u,v}$ that depends on it. Therefore, every bad $(D_0, M_0)$-matched 4-cycle in $S_{u,v}$ different from $R$ has in-degree at least 1 in $\overrightarrow{G}(S_{u,v})$. Therefore, by our earlier observations, every bad $(D_0, M_0)$-matched 4-cycle in $S_{u,v}$ different from $R$ has in-degree exactly 1 in $\overrightarrow{G}(S_{u,v})$.

For Part (c), suppose to the contrary that $\overrightarrow{G}(S_{u,v})$ contains a directed cycle $C:C_1C_2 \ldots C_kC_1$ for some $k \geq 2$. By Part (a), we know that $R$ has in-degree 0 and therefore cannot belong to this cycle. However, there is a (directed) path from $R$ to every other vertex in $\overrightarrow{G}(S_{u,v})$. Among all vertices in the directed cycle $C$, let $C_i$ be chosen so that the distance from $R$ to $C_i$ in $\overrightarrow{G}(S_{u,v})$ is a minimum where $i \in [k]$. Let $P$ be a shortest (directed) path from $R$ to $C_i$ in $\overrightarrow{G}(S_{u,v})$ and let $B$ be the vertex that immediately precedes $C_i$ on the path $P$ (possibly, $B = R$). Since the distance from $R$ to $B$ in $\overrightarrow{G}(S_{u,v})$ is less than the distance from $R$ to $C_i$ in $\overrightarrow{G}(S_{u,v})$, the vertex $B$ does not belong to the directed cycle $C$. Therefore, $C_i$ has in-degree at least 2 in $\overrightarrow{G}(S_{u,v})$, contradicting Part (b). This completes the proof of the claim. (c)

By Claim A.4(c), if there is a cycle in $\overrightarrow{G}(S_{u,v})$, it cannot be an oriented cycle. But then some vertex in that cycle must have in-degree at least 2, contradicting Claim A.4(b). Hence, Claim A.4 implies that $\overrightarrow{G}(S_{u,v})$ is an oriented tree rooted at $R$, and we have proved Property (b) of the statement of Claim A.

It remains to prove Property (c) in the statement of Claim A. Let $C \in S_{u,v}$ be a bad $(D_0, M_0)$-matched 4-cycle where both $u_C'$ and $v_C'$ have an $M_0$-unmatched neighbor, $u''$ and $v''$ say, respectively. These neighbors are clearly distinct, since otherwise $u_C'$ and $v_C'$ are open twins. By Theorem 4 the two $M_0$-unmatched vertices $u''$ and $v''$ belong to distinct odd components of $G - X$. Suppose, to the contrary, that $u_C \in X$. Then, by Theorem 4, $u_C' \notin X$ and $u_C'$ belongs to an odd component of $G - X$ that contains no $M_0$-unmatched vertex. However, $u_C'$ is adjacent to the $M_0$-unmatched vertex $u''$ which implies that $u_C'$ belongs to the same odd component of $G - X$ as the $M_0$-unmatched vertex $u''$, a contradiction. Hence, $u_C \notin X$. Analogously, $v_C \notin X$. If neither $u_C'$ nor $v_C'$ belong to $X$, then $u''$ and $v''$ belong to the same components of $G - X$, a contradiction. Hence, renaming $u_C'$ and $v_C'$, if necessary, we may assume that $v_C \notin X$. Thus, $u_C$ belongs to an odd component of $G - X$ that contains no $M_0$-unmatched vertex. If $u_C' \notin X$, then $v_C$ would be in the same odd component of $G - X$ as the $M_0$-unmatched vertex $u''$, a contradiction. Hence, $u_C \in X$. This establishes Property (c) in the statement of Claim A and completes the proof of Claim A. Q.E.D.
An example of a subgraph of $G$ corresponding to a set $S_{u,v}$ that contains six bad $(D_0, M_0)$-matched 4-cycles is illustrated in Figure 3(a) (where the edges of $M_0$ are thickened, black vertices belong to $D_0$ and white vertices do not, and square white vertices are $M_0$-unmatched vertices).

![Diagram](image)

(a) The structure around the $D_0$-bad pair. 
(b) The oriented tree $\overrightarrow{G}(S_{u,v})$. 
(c) The modification of $D_0$ around $(u_R, v_R)$ when associating it with the leaf $C_5$ of $\overrightarrow{G}(S_{u,v})$. The vertices of each circled pair get swapped in $D_0$.

Figure 3: Example of a $D_0$-bad pair $\{u_R, v_R\}$ with the set of bad $(D_0, M_0)$-matched 4-cycles $S_{u,v} = \{R, C_1, \ldots, C_5\}$. The edges of $M_0$ are thickened; squared vertices are $M_0$-unmatched; black vertices belong to $D_0$.

We now return to the proof of Theorem 2. Our strategy is to modify the set $D_0$ in such a way that the resulting set becomes a locating-dominating set of $G$ of cardinality at most one-half the order of $G$. Consider the set of $D_0$-bad pairs. By Claim A, each such $D_0$-bad pair $\{u, v\}$ belongs to a bad $(D_0, M_0)$-matched 4-cycle $R$ and there is a set $S_{u,v}$ of vertex-disjoint bad $(D_0, M_0)$-matched 4-cycles such that $\overrightarrow{G}(S_{u,v})$ is an oriented tree rooted at $R$. We note that, given two $D_0$-bad pairs $\{u, v\}$ and $\{x, y\}$, the trees $\overrightarrow{G}(S_{u,v})$ and $\overrightarrow{G}(S_{x,y})$ are vertex-disjoint, and furthermore no bad $(D_0, M_0)$-matched 4-cycles of $\overrightarrow{G}(S_{u,v})$ and $\overrightarrow{G}(S_{x,y})$ share any vertex. Indeed, by similar arguments as in Claim A.3 and Claim A.4 we would otherwise contradict the definition of a bad $(D_0, M_0)$-matched 4-cycle.

Now, given a $D_0$-bad pair $\{u, v\}$, consider any leaf $C$ of $\overrightarrow{G}(S_{u,v})$. By Claim A.1(a), the vertices $u'_C$ and $v'_C$ both have a distinct $M_0$-unmatched neighbor, say $u''_C$ and $v''_C$, respectively. Further, by Claim A.1(c), both $u'_C$ and $v'_C$ belong to $X$.

For each $D_0$-bad pair $\{u, v\}$, we select an arbitrary leaf $C$ of $\overrightarrow{G}(S_{u,v})$ and associate the pair of vertices $u'' = u''_C$ and $v'' = v''_C$, of $M_0$-unmatched neighbors of $u'_C$ and $v'_C$, respectively, with the pair $\{u, v\}$, and we write $f(u, v) = \{u'', v''\}$. Let $V^*$ be the set of all $M_0$-unmatched vertices associated with some $D_0$-bad pair. We define the (multi)graph $G^*$ on the vertex set $V^*$ by adding an edge joining $u''$ and $v''$ for each $D_0$-bad pair $\{u, v\}$ such that $f(u, v) = \{u'', v''\}$. As remarked earlier,
the vertices $u''$ and $v''$ are distinct, implying that $G^*$ has no loops (although it may have multiple edges), no isolated vertices, and is subcubic (that is, has maximum degree at most 3). Our aim is to add at most $|V^*|/2$ vertices to $D_0$ and to locally modify $D_0$ around the $D_0$-bad pairs in order to obtain a locating-dominating set, $D'$, of cardinality

$$|D'| \leq \alpha'(G) + \frac{|V^*|}{2} \leq \alpha'(G) + \frac{n - 2\alpha'(G)}{2} = \frac{n}{2}. $$

We now describe the construction of such a set $D'$. Let $D^*$ be a minimum dominating set of $G^*$. Since $G^*$ has no isolated vertex, $|D^*| \leq |V^*|/2$. Since $G^*$ has maximum degree at most 3 and since every vertex outside $D^*$ is adjacent to at least one vertex of $D^*$ in $G^*$, we note that $G^* - D^*$ has maximum degree at most 2. We now build a locating-dominating set from $D_0$ by adding $D^*$ to $D_0$ and by propagating modifications of $D_0$ along the oriented trees associated with all $D_0$-bad pairs. More precisely, we perform our propagation as follows.

**Step 1**: We first consider all $D_0$-bad pairs associated with a pair of vertices of $G^*$ at least one vertex of which belongs to the set $D^*$. Let $\{u, v\}$ be such a $D_0$-bad pair, and let $u''$ and $v''$ be the vertices of $V^*$ such that $f(u, v) = \{u'', v''\}$. Adopting our earlier notation, let $u'' = u''_C$ and $v'' = v''_C$, where $C$ is the chosen leaf in the tree $\vec{G}(S_{u,v})$. Let $R$ be the bad $(D_0, M_0)$-matched 4-cycle in $\vec{S}_{u,v}$ containing $u$ and $v$. Renaming $u''$ and $v''$, if necessary, we may assume that $u''$ belongs to $D^*$. We now consider the unique (directed) path $P$ of $\vec{G}(S_{u,v})$ joining $R$ to $C$ and we modify $D_0$ along $P$ as follows. First, replace $u''_C$ with $u_C$ in $D_0$. If $B$ is the parent of $C$ in $\vec{G}(S_{u,v})$ (and so, $B$ is the vertex on the $(R, C)$-path $P$ that immediately precedes $C$) and $B$ is dependent on $C$ via $x''_B$, where $x_B \in \{u_B, v_B\}$, we replace $x''_B$ with $x_B$ in $D_0$. We continue this process until we perform the modification in the root $R$. This exchange argument in the oriented tree $\vec{G}(S_{u,v})$ associated with the subgraph of $G$ corresponding to the set $S_{u,v}$ illustrated in Figure 3(a) is shown in Figure 3(c). This process is done for all $D_0$-bad pairs associated with a pair of vertices of $G^*$ with at least one member in $D^*$. Let $D'$ be the resulting modified set $D_0$.

**Claim B.** The set of $(D' \cup D^*)$-bad pairs is a proper subset of the set of $D_0$-bad pairs.

**Proof of Claim B.** Let $\{u, v\}$ be an original $D_0$-bad pair associated with a pair of vertices of $G^*$ at least one of which belongs to the set $D^*$. Since at least one of $u$ and $v$ now belongs to $D'$, the pair $\{u, v\}$ is not a $(D' \cup D^*)$-bad pair. It suffices to check the pairs of vertices that could possibly have been affected by the exchange arguments; that is, all vertices previously dominated by a vertex that has been removed from $D_0$ to construct $D'$ (this includes all vertices removed from $D_0$ to construct $D'$). A vertex affected by the modification belongs to a bad $(D_0, M_0)$-matched 4-cycle $A$ in the selected path of $\vec{G}(S_{u,v})$ for some $D_0$-bad pair $\{u, v\}$ that is associated with a pair of vertices of $G^*$ at least one of which belongs to the set $D^*$. Let the vertex set of $A$ be $\{x_A, y_A, x'_A, y'_A\}$ with $\{x_A, y_A\} = \{u_A, v_A\}$, where $x'_A$ has been replaced with $x_A$ in $D'$.

It is sufficient to check that the vertices $x'_A$, $y_A$ and the neighbor, $z$, say, of $x'_A$ not in $A$ are located by $D' \cup D^*$ or belong to $D' \cup D^*$. We observe that even though $D'$ might not belong to $D(M)$, the set $D'$ contains exactly one vertex from each edge in $M_0$. We note that every vertex that was removed from $D_0$ during the exchange arguments when constructing $D'$ is either adjacent to a vertex of $D^*$ or is adjacent to no $M_0$-unmatched vertex. Hence, the vertices of $D_0$ that are adjacent to an $M_0$-unmatched vertex that does not belong to $D^*$ are not removed from $D_0$ during the exchange arguments, implying by Lemma 5 that every $M_0$-unmatched vertex is adjacent to at least two vertices in $D' \cup D^*$ or belongs to $D^*$. It follows that every vertex that is 1-dominated by $D' \cup D^*$ is located by this set. In particular, irrespective of whether $y = u$ or $y = v$, the vertex $y_A$ is 1-dominated by $D' \cup D^*$ and is thus located by $D' \cup D^*$.  

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Suppose that \( z \) is an \( M_0 \)-unmatched vertex. Then, by construction, \( z \in D^* \). In this case, \( x'_{A} \) is dominated by \( z \) and \( x_{A} \) but by no other vertex in \( D' \cup D^* \). If another vertex \( w \) is also only dominated by \( z \) and \( x_{A} \) from \( D' \cup D^* \), then such a vertex cannot be \( M_0 \)-unmatched because the set of \( M_0 \)-unmatched vertices forms an independent set. But then \( w \) is dominated by \( x_{A} \), and its \( M_0 \)-matched neighbor, a contradiction. Hence, if \( z \) is \( M_0 \)-unmatched, then \( x'_{A} \) is located by \( D' \cup D^* \).

Suppose that \( z \) is not an \( M_0 \)-unmatched vertex. Thus, \( z \) belongs to another bad \((D_0, M_0)\)-matched 4-cycle \( B \) of \( \overline{G}(S_{u,v}) \) where \( z = u_B \) and where \( A \) is dependent on \( B \) via \( x'_{A} \) (as illustrated in Figure 2). If \( z \notin D' \), then both \( z \) and \( x'_{A} \) are \( D' \cup D^* \)-dominated and hence are located by \( D' \cup D^* \). Finally, if \( z \in D' \), then \( x_{A} \) is only dominated by \( z \) and \( x_{A} \) from \( D' \cup D^* \). Suppose, to the contrary, that some other vertex \( w \) is also only dominated by \( z \) and \( x_{A} \) from \( D' \cup D^* \). Then, \( w \) must be the neighbor of \( z \) in \( B \) that was removed from \( D_0 \), namely the vertex \( w = u'_B \) (recall that \( z = u_B \)). Thus, \( u'_B \) is adjacent to \( x_{A} \). If \( A = R \), then \( x_{A} \in \{u, v\} \) and \( x_{A} \) would be 3-dominated by \( D_0 \), a contradiction. Hence, \( A \neq R \). If \( x = v \), then we contradict the fact that \( v_{A} \) is adjacent to exactly two vertices of \( D_0 \), namely to \( u'_{A} \) and \( v'_{A} \), and therefore could not be adjacent to \( u'_B \in D \). Hence, \( x = u \). But then \( B \) would be dependent on \( A \) via \( w = u'_B \). However, recall that \( A \) is dependent on \( B \) via \( u'_{A} \), implying that \( \overline{G}(S_{u,v}) \) would contain a 2-cycle joining \( A \) and \( B \), a contradiction. Hence, if \( z \) is not an \( M_0 \)-unmatched vertex, then once again \( x'_{A} \) is located by \( D' \cup D^* \). This completes the proof of Claim [B] (c).

By Claim [B] the set of \( D' \)-bad pairs is a proper subset of the set of \( D_0 \)-bad pairs, implying that all remaining \( D' \)-bad pairs are associated with a pair of vertices of \( G^* \) neither of which belong to the set \( D^* \).

**Step 2:** We next consider all remaining \( D' \)-bad pairs associated with a pair of vertices of \( G^* \) neither of which belong to the set \( D^* \). For each such \( D' \)-bad pair \( \{u, v\} \), we have \( f(u, v) = \{u'', v''\} \) where \( \{u'', v''\} \subseteq V^* \setminus D^* \). Let \( C \) be a component of \( G^* - D^* \) that contains at least one edge. As observed earlier, \( G^* - D^* \) has maximum degree at most 2. Thus, \( C \) is a path or a cycle. If \( C \) is a path, then \( C \) is given by \( c_{0}c_{1} \ldots c_{k-1}c_{0} \) (possibly, \( C \) is a 2-cycle). We now consider an edge \( c_{i}c_{(i+1) \mod k} \) in \( C \), where \( i \in 0, \ldots, k-2 \) if \( C \) is a path and where \( i \in \{0, \ldots, k-1\} \) if \( C \) is a cycle. Let \( \{u, v\} \) be a \( D_0 \)-bad pair such that \( f(u, v) = \{c_{i}, c_{(i+1) \mod k}\} \), and let \( B \) be the bad \((D_0, M_0)\)-matched 4-cycle of \( S_{u,v} \) such that one of \( c_{i} \) and \( c_{(i+1) \mod k} \) is adjacent to \( u'_B \) and the other to \( v'_B \). Let \( c_{i} \) be the neighbor of \( x_{B} \), where \( x_{B} \in \{u_{B}, v_{B}\} \). We now propagate modifications of \( D' \) along a path in \( \overline{G}(S_{u,v}) \) in the same way as we did in Step 1, except that we start the modifications of \( D' \) along the oriented tree by replacing \( x'_{B} \) with \( x_{B} \) and then continuing exactly as before. The resulting modification of \( D' \) ensures that for every vertex in \( C \), at most one of its neighbors is removed from \( D_0 \). This process is done for all \( D_0 \)-bad pairs associated to a pair of vertices of \( G^* \) neither of which belong to the set \( D^* \). Let \( D'' \) be the resulting modified set \( D_0 \).

**Claim C.** No \( D_0 \)-bad pair is a \((D'' \cup D^*)\)-bad pair. Further, the set of \((D'' \cup D^*)\)-bad pairs is a proper subset of the set of \((D' \cup D^*)\)-bad pairs.

**Proof of Claim C.** It suffices for us to check as before the pairs of vertices that could possibly have been affected by the exchange arguments; that is, all vertices previously dominated by a vertex that has been removed from \( D' \) to construct \( D'' \) as well as all vertices removed from \( D' \) to construct \( D'' \). The proof is the same as in the proof of Claim [B] except for the vertices in \( G^* - D^* \). We therefore only prove that vertices that belong to components of \( G^* - D^* \) that contain at least one edge are located by \( D'' \cup D^* \). Let \( c \) be a such a vertex in \( G^* - D^* \). As observed earlier, such a vertex \( c \) belongs to either a path component or a cycle component of \( G^* - D^* \). Further, the modifications of \( D' \) when constructing \( D'' \) ensure that for every vertex in \( G^* - D^* \), at most one of its neighbors
is removed from $D_0$.

We show next that $c$ was 3-dominated by $D_0$. Suppose to the contrary that the vertex $c$ is not 3-dominated by $D_0$ and therefore, by definition of $D(M)$, is adjacent to both ends of some edge $pq$ of $M_0$. In this case, since $c$ has degree at least 1 in $G^* - D^*$ and therefore degree at least 2 in $G^*$, the edge $pq$ must be an edge $x_Bx'_B$, where $x_B \in \{u_B, v_B\}$, in a bad $(D_0, M_0)$-matched 4-cycle $B$ of $S_{u,v}$ for some $D_0$-bad pair $\{u, v\}$. However by Claim A(c), the vertex $x'_B$ belongs to $X$. Therefore the vertex $x_B$, which is $M_0$-matched to $x'_B$, belongs to an odd component of $G - X$ that contains no $M_0$-unmatched vertex. However, the $M_0$-unmatched vertex $c$, which is adjacent to $x_B$, belongs to the same component of $G - X$ as $x_B$, a contradiction. Hence, the vertex $c$ was 3-dominated by $D_0$.

We show now that vertex $c$ is located by $D'' \cup D^*$. If $c$ is 3-dominated by $D'' \cup D^*$, then this follows from the twin-freeness of $G$. Hence we may assume that $c$ is not 3-dominated by $D'' \cup D^*$. Since the vertex $c$ was 3-dominated by $D_0$, and at most one of its neighbors is removed from $D_0$, this implies that the vertex $c$ has exactly two neighbors in $D''$ (and no neighbors in $D^*$ in $G$), and is therefore 2-dominated by $D'' \cup D^*$. Suppose to the contrary that there is a vertex $w$ that is not located from $c$ by $D'' \cup D^*$. Let $d$ be a vertex in $D^*$ that is adjacent to $c$ in $G^*$. Then there exists a $D_0$-bad pair $\{u, v\}$ such that $f(u, v) = \{c, d\}$. Let $B$ be the bad $(D_0, M_0)$-matched 4-cycle of $S_{u,v}$ such that one of $c$ and $d$ is adjacent to $u'_B$ and the other to $v'_B$. Let $c$ be the neighbor of $x'_B$, where $x_B \in \{u_B, v_B\}$. By Step 1, we know that $x'_B \in D''$ and $y_B \in D''$. Therefore, the vertex $w$ must be the vertex $x_B$. Let $z$ be the neighbor of $c$ in $D''$ that is different from $x'_B$. Since the set of neighbors of $c$ in $D''$ is a subset of the set of its neighbors in $D_0$, we note that $\{x'_B, z\} \subset D_0$. If $x_B = v_B$, then $v_B$ would be 3-dominated by $D_0$, contradicting the fact that $B$ is a bad $(D_0, M_0)$-matched 4-cycle. Similarly, if $x_B = u$, then $u$ would be 3-dominated by $D_0$, a contradiction. Hence, $x_B = u_B$ and in $S_{u,v}$, there is a bad $(D_0, M_0)$-matched 4-cycle that depends on $B$ via the vertex $z$. But then $z$ has at least two neighbors apart from $c$ and $x_B$, contradicting the fact that $G$ is cubic. Therefore, the vertex $c$ is located by $D'' \cup D^*$. This completes the proof of Claim C.

Claim C implies that there is no $(D'' \cup D^*)$-bad pair. Thus, the set $D'' \cup D^*$ is a locating-dominating set of $G$. Therefore,

$$
\gamma_L(G) \leq |D| + |D^*| \leq \alpha'(G) + \gamma(G^*) \leq \alpha'(G) + \frac{|V^*|}{2} \leq \alpha'(G) + \frac{n - 2 \alpha'(G)}{2} = \frac{n}{2}
$$

This completes the proof of Theorem 2.

3.3 Tight examples

We remark that the prisms $C_3 \Box K_2$ and $C_4 \Box K_2$ (shown in Figure 4(a) and 4(b), respectively) have location-domination number exactly one-half their order. However, it remains an open problem to characterize all twin-free cubic graphs $G$ of order $n$ that satisfy $\gamma_L(G) = \frac{n}{2}$. Note that the prisms $C_k \Box K_2$ for $k \geq 5$ do not belong to this family.

![Figure 4: The prisms $C_3 \Box K_2$ and $C_4 \Box K_2$.](image-url)
4 Conclusion

We conclude the paper with several intriguing open problems and questions that we have yet to solve.

Problem 1. Characterize the extremal graphs that achieve equality in the bound of Theorem 2; that is, characterize the connected twin-free cubic graphs having location-domination number exactly one-half their order.

Problem 2. Determine whether the result of Theorem 2 can be strengthened by proving Conjecture 1 for subcubic graphs.

Problem 3. Determine whether Theorem 2 can be extended to connected cubic graphs in general (allowing twins) with the exception of a finite set of forbidden graphs. Two such forbidden graphs are the complete graph $K_4$ and the complete bipartite graph $K_{3,3}$, but it is possible that these are the only two exceptions. Proving this would still be weaker than proving the conjecture of Henning and Löwenstein [12] that every cubic graph different from $K_4$ and $K_{3,3}$ has a total locating-dominating set of size at most one-half its order.

Problem 4. Determine whether every connected twin-free cubic graph $G$ satisfies $\gamma_L(G) \leq \alpha'(G)$. More generally, determine classes of twin-free graphs $G$ satisfying $\gamma_L(G) \leq \alpha'(G)$. We remark that Garijo et al. [9] proved that every nontrivial twin-free graph $G$ without 4-cycles satisfies $\gamma_L(G) \leq \alpha'(G)$, and therefore Conjecture 1 holds for these graphs.

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