Quantum Markov chains, sufficiency of quantum channels, and Rényi information measures

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Abstract

A short quantum Markov chain is a tripartite state $\rho_{ABC}$ such that system $A$ can be recovered perfectly by acting on system $C$ of the reduced state $\rho_{BC}$. Such states have conditional mutual information $I(A;B|C)$ equal to zero and are the only states with this property. A quantum channel $\mathcal{N}$ is sufficient for two states $\rho$ and $\sigma$ if there exists a recovery channel using which one can perfectly recover $\rho$ from $\mathcal{N}(\rho)$ and $\sigma$ from $\mathcal{N}(\sigma)$. The relative entropy difference $D(\rho\|\sigma) - D(\mathcal{N}(\rho)\|\mathcal{N}(\sigma))$ is equal to zero if and only if $\mathcal{N}$ is sufficient for $\rho$ and $\sigma$. In this paper, we show that these properties extend to Rényi generalizations of these information measures which were proposed in [Berta et al., J. Math. Phys. 56, 022205, (2015) and arXiv:1410.1443], thus providing an alternate characterization of short quantum Markov chains and sufficient quantum channels. These results give further support to these quantities as being legitimate Rényi generalizations of the conditional mutual information and the relative entropy difference. Along the way, we solve some open questions of Ruskai and Zhang, regarding the trace of particular matrices that arise in the study of monotonicity of relative entropy under quantum operations and strong subadditivity of the von Neumann entropy.

1 Introduction

Markov chains and sufficient statistics are two fundamental notions in probability [Fel97, Nor97] and statistics [Ric94]. Three random variables $X$, $Y$, and $Z$ constitute a three-step Markov chain (denoted as $X - Y - Z$) if $X$ and $Z$ are independent when conditioned on $Y$. In particular, if $p_{XYZ}(x,y,z)$ is their joint probability distribution, then

$$p_{XYZ}(x,y,z) = p_X(x) p_{Y|X}(y|x) p_{Z|Y}(z|y) = p_{X|Y}(x|y) p_{Z|Y}(z|y) p_Y(y).$$  (1.1)

In the information-theoretic framework, such a Markov chain corresponds to a recoverability condition in the following sense. Consider $X$, $Y$, and $Z$ to be the inputs and outputs of two channels (i.e., stochastic maps) $p_{Y|X}$ and $p_{Z|Y}$, as in the figure below. If $X - Y - Z$ is a three-step Markov

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chain, then the input $X$, if lost, can be recovered from $Y$ alone (without any knowledge of $Z$) by
the action of the stochastic map $p_{X|Y}$, as is evident from (1.1).

It is well known that such a Markov chain $X \rightarrow Y \rightarrow Z$ can be characterized by an information
measure \cite{CT91}, namely, the conditional mutual information $I(X; Z|Y)$. For any three random
variables $X$, $Y$ and $Z$, it is defined as

$$I(X; Z|Y) \equiv H(XY) + H(ZY) - H(Y) - H(XYZ),$$

(1.2)

where $H(W) \equiv -\sum_w p_W(w) \log p_W(w)$ is the Shannon entropy of a random variable $W \sim p_W(w)$. It is non-negative and equal to zero if and only if $X \rightarrow Y \rightarrow Z$ is a Markov chain.

In statistics, for a given sample of independent and identically distributed data conditioned on
an unknown parameter $\theta$, a sufficient statistic is a function of the sample whose value contains all the
information needed to compute any estimate of the parameter. One can extend this notion to that
of a sufficient channel (or sufficient stochastic map), as discussed in \cite{Pet86b, Pet88, MP04, Mos05}. A channel $T \equiv T_{X|Y}$ is sufficient for two input distributions $p_X$ and $q_X$ if there exists another channel (a recovery channel) such that both these inputs can be recovered perfectly by sending the outputs of the channel $T_{X|Y}$ corresponding to them through it. This notion of channel sufficiency is likewise characterized by an information measure, namely, the relative entropy difference

$$D(p_X \parallel q_X) - D(T(p_X) \parallel T(q_X)),$$

(1.3)

where $T(p_X)$ and $T(q_X)$ are the distributions obtained after the action of the channel, and $D(p_X \parallel q_X)$ denotes the relative entropy (or Kullback-Leibler divergence) \cite{CT91} between $p_X$ and $q_X$. It is defined as

$$D(p_X \parallel q_X) \equiv \sum_x p_X(x) \log \left( \frac{p_X(x)}{q_X(x)} \right),$$

(1.4)

whenever $q_X(x) \neq 0$ if $p_X(x) \neq 0$ for all $x$, and it is equal to $+\infty$ otherwise. The notion of recoverability provides a connection between the notions of Markov chains and sufficient channels.

The generalization of the above ideas to quantum information theory has been a topic of con-
tinuing and increasing interest (see e.g. \cite{HMP11, BSW15, FR14} and references therein). In the quantum setting, density operators play a role analogous to that of probability distributions in the classical case, and in \cite{HJPW04}, a quantum Markov chain $A \rightarrow C \rightarrow B$ was defined to be a tripartite density operator $\rho_{ABC}$ with conditional (quantum) mutual information $I(A; B|C)_\rho$ equal to zero, where

$$I(A; B|C)_\rho \equiv H(AC)_\rho + H(BC)_\rho - H(C)_\rho - H(ABC)_\rho,$$

(1.5)

and $H(F)_\sigma \equiv -\text{Tr}\{\sigma_F \log \sigma_F\}$ denotes the von Neumann entropy of a density operator $\sigma_F$. (We take the convention $A \rightarrow C \rightarrow B$ for a quantum Markov chain because we are often interested in quantum correlations between Alice ($A$) and Bob ($B$), which are potentially mediated by a third party, here labeled by $C$.) Strong subadditivity of the von Neumann entropy guarantees that the conditional
mutual information $I(A; B|C)_{\rho}$ is non-negative for all density operators [LR73a, LR73b], and it is equal to zero if and only if there is a decomposition of $\mathcal{H}_C$ as

$$\mathcal{H}_C = \bigoplus_j \mathcal{H}_{C_{Lj}} \otimes \mathcal{H}_{C_{Rj}}$$

such that

$$\rho_{ABC} = \bigoplus_j q(j) \rho_{AC_{Lj}} \otimes \rho_{C_{Rj}B},$$

for a probability distribution $\{q(j)\}$ and sets of density operators $\{\rho_{AC_{Lj}}, \rho_{C_{Rj}B}\}$ [HJPW04]. Following [HJPW04], we call such states short quantum Markov chains $A \rightarrow C \rightarrow B$. In analogy with the classical case, $I(A; B|C)_{\rho} = 0$ is equivalent to the full state $\rho_{ABC}$ being recoverable after the loss of system $A$ by the action of a quantum recovery channel $R_{C\rightarrow AC}$ on system $C$ alone:

$$I(A; B|C)_{\rho} = 0 \iff \rho_{ABC} = R_{C\rightarrow AC}(\rho_{BC}),$$

where $R_{C\rightarrow AC}(\cdot) \equiv \frac{1}{2} \rho_{AC}^{1/2} [N(\cdot)]^{-1/2} \rho_{AC}^{1/2}$ is a special case of the so-called Petz recovery channel [Pet86, Pet88].

As a generalization of the classical notion of a sufficient channel, several works have discussed and studied the notion of a sufficient quantum channel [Pet86, Pet88, MP04, Mos05]. Let $\rho$ and $\sigma$ be density operators acting on a Hilbert space $\mathcal{H}$, and let $\mathcal{N}$ be a quantum channel acting on these density operators to density operators acting on a Hilbert space $\mathcal{K}$. Then the quantum channel $\mathcal{N}$ is sufficient for them if one can perfectly recover $\rho$ from $\mathcal{N}(\rho)$ and $\sigma$ from $\mathcal{N}(\sigma)$ by the action of a quantum recovery channel $\mathcal{R}$, i.e., if there exists an $\mathcal{R}$ such that

$$\rho = (\mathcal{R} \circ \mathcal{N})(\rho), \quad \sigma = (\mathcal{R} \circ \mathcal{N})(\sigma).$$

If this condition is true for some recovery channel $\mathcal{R}$, it is known that the following Petz recovery channel $\mathcal{R}_{\sigma, \mathcal{N}}^{P}$ satisfies (19) as well [Pet88]:

$$\mathcal{R}_{\sigma, \mathcal{N}}^{P}(\omega) \equiv \sigma^{1/2} N^{\dagger} \left( [\mathcal{N}(\sigma)]^{-1/2} \omega [\mathcal{N}(\sigma)]^{-1/2} \right) \sigma^{1/2}. \quad (1.10)$$

As a generalization of the classical case, the sufficiency of a quantum channel is characterized by the following information measure, the relative entropy difference

$$D(\rho\|\sigma) - D(\mathcal{N}(\rho)\|\mathcal{N}(\sigma)),$$

where $D(\rho\|\sigma)$ denotes the quantum relative entropy [Ume62]. It is defined as

$$D(\rho\|\sigma) \equiv \text{Tr} \{ \rho [\log \rho - \log \sigma] \},$$

whenever the support of $\rho$ is contained in the support of $\sigma$ and it is equal to $+\infty$ otherwise. The relative entropy difference in (1.11) is non-negative due to the monotonicity of relative entropy under quantum channels [Lin75, Uhl77], and it is equal to zero if and only if $\mathcal{N}$ is sufficient for $\rho$ and $\sigma$, i.e., if the Petz recovery channel $\mathcal{R}_{\sigma, \mathcal{N}}^{P}$ satisfies (19) [Pet86, Pet88]. Further, Mosonyi and Petz have shown that the relative entropy difference in (1.11) is equal to zero if and only if $\rho$, $\sigma$, and $\mathcal{N}$ have the explicit form recalled below in Theorem 6 of Section 3 [MP04, Mos05], which generalizes the result stated in (1.6)-(1.7).
Due to its operational interpretation in the quantum Stein’s lemma [HP91], the quantum relative entropy plays a central role in quantum information theory. In particular, fundamental limits on the performance of information-processing tasks in the so-called “asymptotic, memoryless (or i.i.d.) setting” are given in terms of quantities derived from the quantum relative entropy.

There are, however, other generalized relative entropies (or divergences) which are also of operational significance. Important among these are the Rényi relative entropies [Rény61, Pet86a] and the more recently defined sandwiched Rényi relative entropies [MLDS13, WWY14]. For \( \alpha \in (0, 1) \), the Rényi relative entropies arise in the quantum Chernoff bound [ACMnT+07], which characterizes the probability of error in discriminating two different quantum states in the setting of asymptotically many copies. Moreover, in analogy with the operational interpretation of their classical counterparts, the Rényi relative entropies can be viewed as generalized cutoff rates in quantum binary state discrimination [MH11]. The sandwiched Rényi relative entropies find application in the strong converse domain of a number of settings dealing with hypothesis testing or channel capacity [WWY14, MO15, GW15, TWW14, CMW14, HT14].

This motivates the introduction of Rényi generalizations of the conditional mutual information. Two of these generalizations, defined in [BSW15], are given as follows:

\[
I_\alpha(A; |B|; C)_\rho \equiv \frac{1}{\alpha - 1} \log \text{Tr} \left\{ \rho_{ABC}^{\alpha-1} \rho_{AC}^{1-\alpha} / \rho_{AC}^{1/(\alpha-1)/2} \rho_{BC}^{1-\alpha} / \rho_{BC}^{1/(\alpha-1)/2} \right\}, \tag{1.13}
\]

\[
\tilde{I}_\alpha(A; |B|; C)_\rho \equiv \frac{2\alpha}{\alpha - 1} \log \left\| \rho_{ABC}^{\alpha-1} \rho_{AC}^{1-\alpha} / \rho_{AC}^{1/(\alpha-1)/2} \rho_{BC}^{1-\alpha} / \rho_{BC}^{1/(\alpha-1)/2} \right\|_{2,} \tag{1.14}
\]

where \( \alpha \in (0, 1) \cup (1, \infty) \) denotes the Rényi parameter. (Note that we use the notation \( \| A \|_\alpha \equiv \text{Tr}(\| A^\dagger A \|^{1/\alpha}) \) even for \( \alpha \in (0, 1) \), when it is not a norm.) Both these quantities converge to (1.5) in the limit \( \alpha \to 1 \), they are non-negative, and obey several properties of the conditional mutual information defined in (1.5), as shown in [BSW15]. In [SBW14], the authors proposed some definitions for Rényi generalizations of a relative entropy difference, two of which are as follows:

\[
\Delta_\alpha(\rho, \sigma, \mathcal{N}) \equiv \frac{1}{\alpha - 1} \log \text{Tr} \left\{ \rho^\alpha \sigma^{1-\alpha} \mathcal{N}^{\alpha} \left( \mathcal{N}(\sigma)^{\alpha-1} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{1/\alpha} \right) \sigma^{1/\alpha} \right\}, \tag{1.15}
\]

\[
\tilde{\Delta}_\alpha(\rho, \sigma, \mathcal{N}) \equiv \frac{\alpha}{\alpha - 1} \log \left\| \rho^\alpha \sigma^{1-\alpha} \mathcal{N}^{\alpha} \left( \mathcal{N}(\sigma)^{\alpha-1} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{1/\alpha} \right) \sigma^{1/\alpha} \rho^{1/\alpha} \right\|_{\alpha}. \tag{1.16}
\]

The quantities above converge to (1.11) in the limit \( \alpha \to 1 \) in [SBW14].

The quantities defined in (1.13)-(1.16) can be expressed in terms of Rényi relative entropies, namely the \( \alpha \)-Rényi relative entropy and the \( \alpha \)-sandwiched Rényi relative entropy defined in Section 3.1 [BSW15]. The corresponding expressions for the Rényi generalizations of the conditional mutual information are given in (4.5) and (A.6), respectively, and those for the relative entropy difference are given in (3.53) and (3.54), respectively.

2 Summary of results

As highlighted in the Introduction, an important property of the conditional (quantum) mutual information of a tripartite quantum state is that it is always non-negative and vanishes if and only if the state is a short quantum Markov chain. The relative entropy difference of a pair of quantum states and a quantum channel is also non-negative and vanishes if and only if the channel is sufficient for the pair of states. Consequently, it is reasonable to require that Rényi generalizations of these
information measures are also non-negative and vanish under the same necessary and sufficient conditions as mentioned above.

In this paper, we prove these properties for the quantities $I_\alpha$ and $\tilde{I}_\alpha$, and the quantities $\Delta_\alpha$ and $\tilde{\Delta}_\alpha$ defined in the Introduction. This contributes further evidence that $I_\alpha (A; B| C)_\rho$ and $\tilde{I}_\alpha (A; B| C)_\rho$ are legitimate Rényi generalizations of the conditional mutual information $I (A; B| C)_\rho$ and that $\Delta_\alpha (\rho, \sigma, \mathcal{N})$ and $\tilde{\Delta}_\alpha (\rho, \sigma, \mathcal{N})$ are legitimate Rényi generalizations of the relative entropy difference $D (\rho \| \sigma) - D (\mathcal{N} (\rho) \| \mathcal{N} (\sigma))$. In particular, we prove the following:

1. $I_\alpha (A; B| C)_\rho = 0$ for some $\alpha \in (0, 1) \cup (1, 2)$ if and only if $\rho_{ABC}$ is a short quantum Markov chain with a decomposition as in (1.7).

2. $\tilde{I}_\alpha (A; B| C)_\rho = 0$ for some $\alpha \in (1/2, 1) \cup (1, \infty)$ if and only if $\rho_{ABC}$ is a short quantum Markov chain with a decomposition as in (1.7).

3. $\Delta_\alpha (\rho, \sigma, \mathcal{N})$ is non-negative for $\alpha \in (0, 1) \cup (1, 2)$ and $\tilde{\Delta}_\alpha (\rho, \sigma, \mathcal{N})$ is non-negative for $\alpha \in (1/2, 1) \cup (1, \infty)$.

4. $\Delta_\alpha (\rho, \sigma, \mathcal{N}) = 0$ for some $\alpha \in (0, 1) \cup (1, 2)$ if and only if the quantum channel $\mathcal{N}$ is sufficient for states $\rho$ and $\sigma$, so that $\mathcal{N}$, $\rho$, and $\sigma$ decompose as in (3.19)-(3.20).

5. $\tilde{\Delta}_\alpha (\rho, \sigma, \mathcal{N}) = 0$ for some $\alpha \in (1/2, 1) \cup (1, \infty)$ if and only if the quantum channel $\mathcal{N}$ is sufficient for states $\rho$ and $\sigma$, so that $\mathcal{N}$, $\rho$, and $\sigma$ decompose as in (3.19)-(3.20).

6. Generalizations of the conditional mutual information and the relative entropy difference arising from the so-called min- and max-relative entropies [Dat09, DKF+12] (which play an important role in one-shot information theory) satisfy identical properties.

Along the way, we resolve some open questions stated in [Rus02, Zha14]. Let $\rho_{ABC}$ be a positive definite density operator. We prove that

$$\text{Tr} \left\{ \left( \frac{(1-\alpha)^{2\alpha}}{\rho_{AC}} \frac{(1-\alpha)^{1-\alpha}}{\rho_C} \frac{1-\alpha}{\rho_{BC}} \frac{1-\alpha}{\rho_{AC}} \right)^{1/(1-\alpha)} \right\} \leq 1, \quad (2.1)$$

for all $\alpha \in (0, 1) \cup (1, 2)$, and

$$\text{Tr} \left\{ \left( \frac{(1-\alpha)^{2\alpha}}{\rho_{AC}} \frac{(1-\alpha)^{1-\alpha}}{\rho_C} \frac{1-\alpha}{\rho_{BC}} \frac{1-\alpha}{\rho_{AC}} \right)^{(1-\alpha)/(1-\alpha)} \right\} \leq 1, \quad (2.2)$$

for all $\alpha \in (1/2, 1) \cup (1, \infty)$. Let $\rho$ and $\sigma$ be positive definite density operators and let $\mathcal{N}$ be a strict completely positive trace preserving (CPTP) map (that is, a CPTP map such that $\mathcal{N}(X)$ is positive definite whenever $X$ is positive definite). We prove that

$$\text{Tr} \left\{ \left[ \sigma^{(1-\alpha)/2} \mathcal{N}^\dagger \left( \mathcal{N} (\sigma)^{(1-\alpha)/2} \mathcal{N} (\rho)^{1-\alpha} \mathcal{N} (\sigma)^{(1-\alpha)/2} \right)^{(1-\alpha)/2} \sigma^{(1-\alpha)/2} \right]^{1/(1-\alpha)} \right\} \leq 1, \quad (2.3)$$

for $\alpha \in (0, 1)$, and

$$\text{Tr} \left\{ \left[ \sigma^{(1-\alpha)/2} \mathcal{N}^\dagger \left( \mathcal{N} (\sigma)^{(1-\alpha)/2} \mathcal{N} (\rho)^{1-\alpha} \mathcal{N} (\sigma)^{(1-\alpha)/2} \right)^{(1-\alpha)/2} \sigma^{(1-\alpha)/2} \right]^{(1-\alpha)/(1-\alpha)} \right\} \leq 1, \quad (2.4)$$
for $\alpha \in (1/2, 1)$. Determining whether (2.3) holds for $\alpha \in (1, 2)$ and (2.4) for $\alpha \in (1, \infty)$, remains an open problem. By taking the limit $\alpha \rightarrow 1$, the inequalities in (2.3) and (2.4) imply that

$$\text{Tr} \left\{ \exp \left\{ \log \sigma + \mathcal{N}^\dagger (\log \mathcal{N}(\rho) - \log \mathcal{N}(\sigma)) \right\} \right\} \leq 1. \quad (2.5)$$

The rest of the paper is devoted to establishing these claims. We begin by recalling some mathematical preliminaries and known results, and follow by establishing the latter claims first and then move on to the former ones.

## 3 Preliminaries

Let $\mathcal{B}(\mathcal{H})$ denote the algebra of bounded linear operators acting on a Hilbert space $\mathcal{H}$. We restrict to finite-dimensional Hilbert spaces throughout this paper. We denote the support of an operator $A$ by $\text{supp}(A)$. For a Hermitian operator $A$, by $A^{-1}$ we mean the inverse restricted to $\text{supp}(A)$, so that $AA^{-1} = A^{-1}A$ is the orthogonal projection onto $\text{supp}(A)$. More generally, for a function $f$ and Hermitian operator $A$ with spectral decomposition $A = \sum_i \lambda_i |i\rangle \langle i|$, we define $f(A)$ to be $\sum_{i: \lambda_i \neq 0} f(\lambda_i) |i\rangle \langle i|$. Let $\mathcal{B}(\mathcal{H})_+$ denote the subset of positive definite operators, and let $\mathcal{B}(\mathcal{H})_{++}$ denote the subset of positive definite operators. We also write $X \geq 0$ if $X \in \mathcal{B}(\mathcal{H})_+$ and $X > 0$ if $X \in \mathcal{B}(\mathcal{H})_{++}$. An operator $\rho$ is in the set $\mathcal{S}(\mathcal{H})$ of density operators (or states) if $\rho \in \mathcal{B}(\mathcal{H})_+$ and $\text{Tr}\{\rho\} = 1$, and an operator $\rho$ is in the set $\mathcal{S}(\mathcal{H})_{++}$ of positive definite density operators if $\rho \in \mathcal{B}(\mathcal{H})_{++}$ and $\text{Tr}\{\rho\} = 1$. Throughout much of the paper, for technical convenience and simplicity, we consider states in $\mathcal{S}(\mathcal{H})_{++}$. For $\alpha \geq 1$, we define the $\alpha$-norm of an operator $X$ as

$$\|X\|_\alpha \equiv \left[ \text{Tr}\{ (\sqrt{X^\dagger X})^\alpha \} \right]^{1/\alpha}, \quad (3.1)$$

and we use the same notation even for the case $\alpha \in (0, 1)$, when it is not a norm. The fidelity $F(\rho, \sigma)$ of two states $\rho, \sigma \in \mathcal{S}(\mathcal{H})$ is defined as

$$F(\rho, \sigma) \equiv \| \sqrt{\rho} \sqrt{\sigma} \|_1^2. \quad (3.2)$$

A quantum channel (or quantum operation) is given by a completely positive, trace-preserving (CPTP) map $\mathcal{N} : \mathcal{B}(\mathcal{H}) \mapsto \mathcal{B}(\mathcal{K})$, with $\mathcal{H}$ and $\mathcal{K}$ being the input and output Hilbert spaces of the channel, respectively. Let $\langle C, D \rangle \equiv \text{Tr}\{ C^\dagger D \}$ denote the Hilbert-Schmidt inner product of $C, D \in \mathcal{B}(\mathcal{H})$. The adjoint of the quantum channel $\mathcal{N}$ is a completely positive unital map $\mathcal{N}^\dagger : \mathcal{B}(\mathcal{K}) \mapsto \mathcal{B}(\mathcal{H})$ defined through the following relation:

$$\langle B, \mathcal{N}(A) \rangle = \langle \mathcal{N}^\dagger(B), A \rangle, \quad (3.3)$$

for all $A \in \mathcal{B}(\mathcal{H})$ and $B \in \mathcal{B}(\mathcal{K})$. A linear map is said to be a strict CPTP map if it is CPTP and if $\mathcal{N}(A) \in \mathcal{B}(\mathcal{K})_{++}$ for all $A \in \mathcal{B}(\mathcal{H})_{++}$. We denote the identity channel as $\text{id}$ but often suppress it for notational simplicity.

The set $\{ U^i \}$ of Heisenberg-Weyl unitaries acting on a finite-dimensional Hilbert space $\mathcal{H}$ of dimension $d$ has the property that

$$\frac{1}{d^2} \sum_i U^i X U^i \dagger = \text{Tr}\{ X \} \frac{I}{d}, \quad (3.4)$$

for any operator $X$ acting on $\mathcal{H}$. 

6
3.1 Generalized relative entropies

The following relative entropies of a density operator $\rho$ with respect to a positive semidefinite operator $\sigma$ play an important role in this paper. In what follows, we restrict the definitions to the case in which $\rho$ and $\sigma$ satisfy the condition $\text{supp}\, \rho \subseteq \text{supp}\, \sigma$, with the understanding that they are equal to $+\infty$ otherwise. The $\alpha$-Rényi relative entropy is defined for $\alpha \in (0, 1) \cup (1, \infty)$ as follows [Pet86a]:

$$D_\alpha(\rho\|\sigma) \equiv \frac{1}{\alpha - 1} \log \text{Tr} \left\{ \rho^\alpha \sigma^{1-\alpha} \right\}. \quad (3.5)$$

The $\alpha$-sandwiched Rényi relative entropy is defined for $\alpha \in (0, 1) \cup (1, \infty)$ as follows [MLDS+13, WWY14]:

$$\tilde{D}_\alpha(\rho\|\sigma) \equiv \frac{1}{\alpha - 1} \log \text{Tr} \left\{ \left( \sigma^{(1-\alpha)/2\alpha} \rho \sigma^{(1-\alpha)/2\alpha} \right)^\alpha \right\}. \quad (3.6)$$

Both quantities above reduce to the quantum relative entropy in (1.12) in the limit $\alpha \to 1$.

A fundamental property of the quantum relative entropy is that it is monotone under quantum operations (also known as the data-processing inequality):

$$D(\rho\|\sigma) \geq D(\mathcal{N}(\rho)\|\mathcal{N}(\sigma)), \quad (3.7)$$

where $\mathcal{N}$ is a quantum operation. It is known that the data-processing inequality is satisfied by the $\alpha$-Rényi relative entropy for $\alpha \in [0, 1) \cup (1, \infty]$ [Pet86a], and for the $\alpha$-sandwiched Rényi relative entropy for $\alpha \in [1/2, 1) \cup (1, \infty]$ [FL13, Bei13, MLDS+13, WWY14, MO15].

Two special cases of the $\alpha$-sandwiched Rényi relative entropy are of particular significance in one-shot information theory [Ren05, Tom12], namely the min-relative entropy [DKF+12] and the max-relative entropy [Dat09]. These are defined as follows:

$$D_{\text{min}}(\rho\|\sigma) \equiv \tilde{D}_{1/2}(\rho\|\sigma) = -\log F(\rho, \sigma), \quad (3.8)$$

and

$$D_{\text{max}}(\rho\|\sigma) \equiv \inf \{\lambda : \rho \leq 2^\lambda \sigma\} = \lim_{\alpha \to \infty} \tilde{D}_\alpha(\rho\|\sigma). \quad (3.9)$$

The relative entropies defined above, satisfy the following lemma [MLDS+13]:

**Lemma 1** For $\omega \in \mathcal{S}(\mathcal{H})$ and $\tau \in \mathcal{B}(\mathcal{H})_+$, such that $\text{Tr}\{\omega\} \geq \text{Tr}\{\tau\}$,

$$D_\alpha(\omega\|\tau) \geq 0 \text{ for } \alpha \in (0, 1) \cup (1, 2), \quad (3.10)$$

$$\tilde{D}_\alpha(\omega\|\tau) \geq 0 \text{ for } \alpha \in (1/2, 1) \cup (1, \infty), \quad (3.11)$$

$$D_{\text{min}}(\omega\|\tau) \geq 0, \quad (3.12)$$

$$D_{\text{max}}(\omega\|\tau) \geq 0, \quad (3.13)$$

with equalities holding if and only if $\omega = \tau$.

In proving our results, we also employ the notion of a quantum $f$-divergence, first introduced by Petz [Pet86a]. It can be conveniently expressed as follows [TCR09]:
Definition 2 For $A \in \mathcal{B}(\mathcal{H})_+$, $B \in \mathcal{B}(\mathcal{H})_+$, and an operator convex function $f$ on $[0, \infty)$, the $f$-divergence of $A$ with respect to $B$ is given by

$$S_f(A \| B) = \langle \Gamma \mid (\sqrt{B} \otimes I) f (B^{-1} \otimes A^T) (\sqrt{B} \otimes I) \mid \Gamma \rangle,$$

(3.14)

where $|\Gamma\rangle = \sum_i |i\rangle \otimes |i\rangle$, and $\{|i\rangle\}$ is an orthonormal basis of $\mathcal{H}$ with respect to which the transpose is defined.

Remark 3 Special cases of this are as follows:

1. The trace expression $\text{Tr} \{ \rho^\alpha \sigma^{1-\alpha} \}$ of the $\alpha$-Rényi relative entropy, for the choice $f(x) = x^\alpha$ for $\alpha \in (1, 2)$, and $-\text{Tr} \{ \rho^\alpha \sigma^{1-\alpha} \}$ for the choice $f(x) = -x^\alpha$ for $\alpha \in (0, 1)$.

2. The quantum relative entropy for the choice $f(x) = x \log x$.

The equivalence relations given in Lemma 4 below follow directly from [HMPB11, Theorem 5.1].

Lemma 4 For $A, B \in \mathcal{B}(\mathcal{H})_+$ and a strict CPTP map $\mathcal{N}$ acting on $\mathcal{B}(\mathcal{H})$, the following conditions are equivalent:

$$S_f (\mathcal{N}(A) \| \mathcal{N}(B)) = S_f (A \| B) \quad \text{for some operator convex function } f \text{ on } [0, \infty),$$

(3.15)

$$\mathcal{N}^\dagger [\log \mathcal{N}(A) - \log \mathcal{N}(B)] = \log A - \log B.$$  

(3.16)

(One requires some technical conditions for $f$ which are specified in [HMPB11, Theorem 5.1]. These conditions are satisfied by the functions given in Remark 3.)

Lemma 5 If $\rho_{ABC}$ is a positive definite density operator such that

$$\log \rho_{ABC} = \log \rho_{AC} + \log \rho_{BC} - \log \rho_C,$$

(3.17)

then it is a short quantum Markov chain $A \rightarrow C \rightarrow B$.

Proof. The identity (3.17) is known from [Rus02] to be a condition for the conditional mutual information $I(A; B|C)_\rho$ of $\rho_{ABC}$ to be equal to zero, which, by [HJPW04], implies that $\rho_{ABC}$ is a short quantum Markov chain $A \rightarrow C \rightarrow B$. 

Theorem 6 ([MP04, Mos05]) Let $\rho \in S(\mathcal{H})_+$, $\sigma \in \mathcal{B}(\mathcal{H})_+$, and $\mathcal{N}$ be a strict CPTP map. Then $\mathcal{N}$ is sufficient for $\rho$ and $\sigma$ (as in (1.9)) if and only if the following conditions hold

1. There exist decompositions of $\mathcal{H}$ and $\mathcal{K}$ as follows:

$$\mathcal{H} = \bigoplus_j \mathcal{H}_{L_j} \otimes \mathcal{H}_{R_j}, \quad \mathcal{K} = \bigoplus_j \mathcal{K}_{L_j} \otimes \mathcal{K}_{R_j},$$

(3.18)

where $\dim(\mathcal{H}_{L_j}) = \dim(\mathcal{K}_{L_j})$ for all $j$. 

8
2. With respect to the decomposition in (3.18), \( \rho \) and \( \sigma \) can be written as follows:

\[
\rho = \bigoplus_j p(j) \rho^j_L \otimes \tau^j_R, \quad \sigma = \bigoplus_j q(j) \sigma^j_L \otimes \tau^j_R,
\]

for some probability distribution \( \{p(j)\} \), positive reals \( \{q(j)\} \), sets of states \( \{\rho^j_L\} \) and \( \{\tau^j_R\} \) and set of positive definite operators \( \{\sigma^j_L\} \).

3. With respect to the decomposition in (3.18), the quantum channel \( \mathcal{N} \) can be written as

\[
\mathcal{N} = \bigoplus_j \mathcal{U}_j \otimes \mathcal{N}_j^R,
\]

where \( \{\mathcal{U}_j : \mathcal{B}(\mathcal{H}_L_j) \to \mathcal{B}(\mathcal{K}_L_j)\} \) is a set of unitary channels and \( \{\mathcal{N}_j^R : \mathcal{B}(\mathcal{H}_R_j) \to \mathcal{B}(\mathcal{K}_R_j)\} \) is a set of quantum channels. Furthermore, with respect to the decomposition in (3.18), the adjoint of \( \mathcal{N} \) acts as

\[
\mathcal{N}^\dagger = \bigoplus_j \mathcal{U}_j^\dagger \otimes (\mathcal{N}_j^R)^\dagger.
\]

### 3.2 Trace inequalities

The following lemma is a consequence of [Eps73] (see also [CL08, Theorem 1.1]):

**Lemma 7** For \( A \in \mathcal{B}(\mathcal{H}), B \in \mathcal{B}(\mathcal{H})_+ \), and \( p \in (0, 1) \), the map \( B \mapsto \text{Tr}\{ (AB^p A^\dagger)^{1/p} \} \) is concave.

For invertible \( A \in \mathcal{B}(\mathcal{H}), B \in \mathcal{B}(\mathcal{H})_+ \), and \( p \in (-1, 0) \), the map \( B \mapsto \text{Tr}\{ (AB^p A^{\dagger})^{1/p} \} \) is concave.

**Lemma 8** Let \( \rho \in \mathcal{S}(\mathcal{H}), \sigma \in \mathcal{B}(\mathcal{H})_+ \), and let \( \mathcal{N} : \mathcal{B}(\mathcal{H}) \to \mathcal{B}(\mathcal{K}) \) be a CPTP map. For \( \alpha \in (0, 1) \)

\[
\text{Tr} \left\{ \left[ \sigma^{(1-\alpha)/2} \mathcal{N}^\dagger \left( \mathcal{N}(\sigma)^{(\alpha-1)/2} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2} \right)^{1/(1-\alpha)} \right]^{1/(1-\alpha)} \right\} \leq 1. \tag{3.22}
\]

For \( \alpha \in (1/2, 1) \)

\[
\text{Tr} \left\{ \left[ \sigma^{(1-\alpha)/2\alpha} \mathcal{N}^\dagger \left( \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \mathcal{N}(\rho)^{(1-\alpha)/\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \right)^{(1-\alpha)/\alpha} \right]^{\alpha/(1-\alpha)} \right\} \leq 1. \tag{3.23}
\]

**Proof.** We begin by proving (3.22) for \( \alpha \in (0, 1) \). By Stinespring’s dilation theorem [Sti55], a given quantum channel \( \mathcal{N} \) can be realized as

\[
\mathcal{N}(\omega) = \text{Tr}_{E^\dagger} \left\{ U (\omega \otimes |0\rangle_E \langle 0|_E) U^\dagger \right\} \quad \forall \omega \in \mathcal{B}(\mathcal{H}), \tag{3.24}
\]

for some unitary \( U \) taking \( \mathcal{H} \otimes \mathcal{H}_E \) to \( \mathcal{K} \otimes \mathcal{H}_E^\dagger \) and fixed state \( |0\rangle_E \) in an auxiliary Hilbert space \( \mathcal{H}_E \). Furthermore, it suffices to take \( \dim(\mathcal{H}_E) \leq \dim(\mathcal{H}) \dim(\mathcal{K}) \) because the number of Kraus operators for the channel \( \mathcal{N} \) can always be taken less than or equal to \( \dim(\mathcal{H}) \dim(\mathcal{K}) \) and this is the dimension needed for an environment system \( \mathcal{H}_E \). The adjoint of \( \mathcal{N} \) is given by

\[
\mathcal{N}^\dagger(\tau) = \langle 0|_E U^\dagger (\tau \otimes I_{E^\dagger}) U |0\rangle_E \quad \forall \tau \in \mathcal{B}(\mathcal{K}), \tag{3.25}
\]
Then

\[
\begin{align*}
\text{Tr} \left\{ \left( \sigma^{(1-\alpha)/2} N^\dagger \left( N (\sigma)^{(\alpha-1)/2} N (\rho)^{1-\alpha} N (\sigma)^{(\alpha-1)/2} \right) \sigma^{(1-\alpha)/2} \right)^{1/(1-\alpha)} \right\} \\
= \text{Tr} \left\{ \left( \sigma^{(1-\alpha)/2} (0|_E U^\dagger \left( \left[ \left( N (\sigma)^{(\alpha-1)/2} N (\rho)^{1-\alpha} N (\sigma)^{(\alpha-1)/2} \right) \right] \otimes I_{E'} \right) U |0\rangle_E \sigma^{(1-\alpha)/2} \right)^{1/(1-\alpha)} \right\} \\
= \text{Tr} \left\{ \left( AA^\dagger \right)^{1/(1-\alpha)} \right\} \\
\end{align*}
\]

(3.26)

where

\[
A = \sigma^{(1-\alpha)/2} (0|_E U^\dagger K_\alpha^\dagger, \quad (3.27)
K_\alpha \equiv K_\alpha (\rho, \sigma, N) \equiv N (\rho)^{1-\alpha} N (\sigma)^{\alpha-1} \otimes I_{E'}.
(3.28)
\]

Then the above is equal to

\[
\text{Tr} \left\{ \left( A^\dagger A \right)^{1/(1-\alpha)} \right\} = \text{Tr} \left\{ \left( K_\alpha U |0\rangle_E \sigma^{\frac{1-\alpha}{2}} \sigma^{\frac{\alpha}{2}} (0|_E U^\dagger K_\alpha^\dagger) \right)^{1/(1-\alpha)} \right\},
(3.29)
\]

because the eigenvalues of $AA^\dagger$ and $A^\dagger A$ are the same for any operator $A$. Using that

\[
\begin{align*}
U |0\rangle_E \sigma^{(1-\alpha)/2} \sigma^{(1-\alpha)/2} (0|_E U^\dagger &= U \left( \sigma^{1-\alpha} \otimes |0\rangle \langle 0|_E \right) U^\dagger \\
&= U \left( \sigma \otimes |0\rangle \langle 0|_E \right)^{1-\alpha} U^\dagger \\
&= \left[ U \left( \sigma \otimes |0\rangle \langle 0|_E \right) U^\dagger \right]^{1-\alpha},
(3.30)
\end{align*}
\]

the right hand side of (3.29) is equal to

\[
\text{Tr} \left\{ \left( K_\alpha \left[ U \left( \sigma \otimes |0\rangle \langle 0|_E \right) U^\dagger \right]^{1-\alpha} K_\alpha^\dagger \right)^{1/(1-\alpha)} \right\}.
(3.31)
\]

Let \( \{ U_{E'}^i \} \) be a set of Heisenberg-Weyl operators for the $E'$ system and let $\pi_{E'}$ denote the maximally mixed state on system $E'$. Now we use Lemma 7 to establish the inequality below:

\[
\begin{align*}
\text{Tr} \left\{ \left( K_\alpha \left[ U \left( \sigma \otimes |0\rangle \langle 0|_E \right) U^\dagger \right]^{1-\alpha} K_\alpha^\dagger \right)^{1/(1-\alpha)} \right\} \\
= \frac{1}{d_{E'}} \sum_i \text{Tr} \left\{ \left( K_\alpha \left[ U_{E'}^i U \left( \sigma \otimes |0\rangle \langle 0|_E \right) U^\dagger \left( U_{E'}^i \right)^\dagger \right]^{1-\alpha} K_\alpha^\dagger \right)^{1/(1-\alpha)} \right\} \\
\leq \text{Tr} \left\{ \left( K_\alpha \left[ \frac{1}{d_{E'}} \sum_i U_{E'}^i U \left( \sigma \otimes |0\rangle \langle 0|_E \right) U^\dagger \left( U_{E'}^i \right)^\dagger \right]^{1-\alpha} K_\alpha^\dagger \right)^{1/(1-\alpha)} \right\} \\
= \text{Tr} \left\{ \left( K_\alpha \left[ \text{Tr}_{E'} \left\{ U \left( \sigma \otimes |0\rangle \langle 0|_E \right) U^\dagger \otimes \pi_{E'} \right\} \right]^{1-\alpha} K_\alpha^\dagger \right)^{1/(1-\alpha)} \right\}.
\end{align*}
\]

(3.32)
(3.33)
(3.34)
Continuing, the last line above is equal to
\[
\text{Tr} \left\{ \left( K_\alpha \left[ \mathcal{N} \left( \sigma \right) \otimes \pi_{E'} \right] \right)^{1-\alpha} K_\alpha^\dagger \right\}^{1/(1-\alpha)}
\]
\[
= \text{Tr} \left\{ \left( K_\alpha \left[ \mathcal{N} \left( \sigma \right) \otimes \pi_{E'}^{1-\alpha} \right] K_\alpha \right)^{1/(1-\alpha)} \right\}
\]
\[
= \text{Tr} \left\{ \left( \mathcal{N} \left( \rho \right)^{1-\alpha/2} \mathcal{N} \left( \sigma \right)^{(\alpha-1)/2} \mathcal{N} \left( \sigma \right)^{(\alpha-1)/2} \mathcal{N} \left( \rho \right)^{1-\alpha/2} \otimes \pi_{E'}^{1-\alpha} \right)^{1/(1-\alpha)} \right\}
\]
\[
= \text{Tr} \left\{ \left( \mathcal{N} \left( \rho \right)^{1-\alpha/2} \mathcal{N} \left( \sigma \right)^0 \mathcal{N} \left( \rho \right)^{1-\alpha/2} \otimes \pi_{E'}^{1-\alpha} \right)^{1/(1-\alpha)} \right\}
\]
\[
\leq \text{Tr} \left\{ \left( \mathcal{N} \left( \rho \right)^{1-\alpha} \otimes \pi_{E'} \right)^{1/(1-\alpha)} \right\}
\]
\[
= 1,
\]
where the inequality follows because
\[
\mathcal{N} \left( \rho \right)^{1-\alpha/2} \mathcal{N} \left( \sigma \right)^0 \mathcal{N} \left( \rho \right)^{1-\alpha/2} \leq \mathcal{N} \left( \rho \right)^{1-\alpha},
\]
and \( \text{Tr} \{ f (A) \} \leq \text{Tr} \{ f (B) \} \) when \( A \leq B \) and \( f (x) \equiv x^{1/(1-\alpha)} \) is a monotone non-decreasing function on \([0, \infty)\). The other inequality \( (3.39) \) follows from \( (3.22) \) by making the substitution \( \alpha \rightarrow (2\alpha - 1)/\alpha \). ■

As a corollary, we establish the following trace inequality, which was left as an open question in [Rus02]:

**Corollary 9** For \( \rho \in S(\mathcal{H})_{++} \), \( \sigma \in B(\mathcal{H})_{++} \), and a strict CPTP map \( \mathcal{N} \) acting on \( S(\mathcal{H}) \), the following inequality holds:

\[
\text{Tr} \left\{ \exp \left\{ \log \sigma + \mathcal{N}^\dagger \left( \log \mathcal{N} \left( \rho \right) - \log \mathcal{N} \left( \sigma \right) \right) \right\} \right\} \leq 1.
\]

**Proof.** This follows by taking the limit \( \alpha \uparrow 1 \) in Lemma 8 and using [SBW14] Lemma 24:

\[
\exp \left\{ \log \sigma + \mathcal{N}^\dagger \left( \log \mathcal{N} \left( \rho \right) - \log \mathcal{N} \left( \sigma \right) \right) \right\}
= \lim_{\alpha \uparrow 1} \left[ \sigma^{(1-\alpha)/2} \mathcal{N}^\dagger \left( \mathcal{N} \left( \sigma \right)^{(\alpha-1)/2} \mathcal{N} \left( \rho \right)^{(1-\alpha)/2} \mathcal{N} \left( \sigma \right)^{(\alpha-1)/2} \sigma^{(1-\alpha)/2} \right)^{1/(1-\alpha)} \right].
\]

The inequality in Lemma 8 is preserved in the limit. ■

We can also solve an open question from [Zha14] for some values of \( \alpha \):

**Corollary 10** Let \( \rho_{ABC} \in S(\mathcal{H}_{ABC}) \). If \( \alpha \in (0, 1) \), then

\[
\text{Tr} \left\{ \left( \rho_{AC}^{(1-\alpha)/2} \rho_{BC}^{(\alpha-1)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{BC}^{(1-\alpha)/2} \right)^{1/(1-\alpha)} \right\} \leq 1.
\]

If \( \alpha \in (1/2, 1) \), then

\[
\text{Tr} \left\{ \left( \rho_{AC}^{(1-\alpha)/2} \rho_{BC}^{(\alpha-1)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{BC}^{(1-\alpha)/2} \right)^{\alpha/(1-\alpha)} \right\} \leq 1.
\]

If \( \rho_{ABC} \in S(\mathcal{H}_{ABC})_{++} \), then (3.43) holds for \( \alpha \in (1, 2) \) and (3.44) holds for \( \alpha \in (1, \infty) \).
Proof. The inequalities in (3.43)-(3.44) follow directly from Lemma [S] by choosing \( \rho = \rho_{ABC} \), \( \sigma = \rho_{AC} \), and \( \mathcal{N} = \text{Tr}_A \). The fact that (3.43) holds for \( \alpha \in (1, 2) \) and \( \rho_{ABC} \in \mathcal{S}(\mathcal{H}_{ABC})_{++} \) follows by letting \( \beta \in (0, 1) \) be such that \( \alpha + \beta = 2 \) and noticing that

\[
\left( \rho_{AC}^{(1-\alpha)/2} \rho_C^{(\alpha-1)/2} \rho_{AC}^{-1} \rho_{AC}^{1-\alpha} \rho_C^{(\alpha-1)/2} \right)^{1/(1-\alpha)} = \left( \rho_{AC}^{(1-\beta)/2} \rho_C^{(\beta-1)/2} \rho_{AC}^{-1} \rho_{AC}^{1-\beta} \rho_C^{(\beta-1)/2} \right)^{1/(1-\beta)}
\]

when \( \rho_{ABC} \) is positive definite. By a similar line of reasoning, (3.44) holds for \( \alpha \in (1, \infty) \) and \( \rho_{ABC} \in \mathcal{S}(\mathcal{H}_{ABC})_{++} \) by letting \( \beta \in (1/2, 1) \) be such that \( 1/\alpha + 1/\beta = 2 \). ■

Remark 11 Let \( \rho_{ABC} \in \mathcal{S}(\mathcal{H}_{ABC})_{++} \). The well-known inequality [LR73b]

\[
\text{Tr}\{\exp\{\log \rho_{AC} + \log \rho_{BC} - \log \rho_C\}\} \leq 1
\]

(3.46)

follows from Corollary [10] by taking the limit \( \alpha \to 1 \) and using the generalized Lie-Trotter product formula [Suz83]:

\[
\exp\{\log \rho_{AC} + \log \rho_{BC} - \log \rho_C\} = \lim_{\alpha \to 1} \left( \rho_{AC}^{(1-\alpha)/2} \rho_C^{(\alpha-1)/2} \rho_{BC} \rho_C^{(\alpha-1)/2} \rho_{AC}^{(1-\alpha)/2} \right)^{1/(1-\alpha)}.
\]

(3.47)

The following proposition establishes some important properties of the \( \Delta_{\alpha}(\rho, \sigma, \mathcal{N}) \) and \( \bar{\Delta}_{\alpha}(\rho, \sigma, \mathcal{N}) \) quantities, which were left as open questions from [SBW14]:

Proposition 12 Let \( \rho \in \mathcal{S}(\mathcal{H}) \), \( \sigma \in \mathcal{B}(\mathcal{H})_{++} \), and let \( \mathcal{N} \) be a CPTP map. For all \( \alpha \in (0, 1) \)

\[
\Delta_{\alpha}(\rho, \sigma, \mathcal{N}) \geq 0,
\]

(3.48)

with equality occurring if and only if

\[
\rho = \left[ \sigma^{(1-\alpha)/2} \mathcal{N} \left( \mathcal{N}(\sigma)^{(\alpha-1)/2} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2} \right)^{1/(1-\alpha)} \right]^{1/(1-\alpha)}.
\]

(3.49)

For all \( \alpha \in (1/2, 1) \)

\[
\bar{\Delta}_{\alpha}(\rho, \sigma, \mathcal{N}) \geq 0,
\]

(3.50)

with equality occurring if and only if

\[
\rho = \left[ \sigma^{(1-\alpha)/2} \mathcal{N} \left( \mathcal{N}(\sigma)^{(\alpha-1)/2} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2} \right)^{1/(1-\alpha)} \right]^{\alpha/(1-\alpha)}.
\]

(3.51)

Let \( \rho \in \mathcal{S}(\mathcal{H})_{++} \), \( \sigma \in \mathcal{B}(\mathcal{H})_{++} \), and let \( \mathcal{N} \) be a strict CPTP map. Then (3.48) holds for \( \alpha \in (1, 2) \) and equality occurs if and only if

\[
\left[ \mathcal{N}(\sigma)^{(\alpha-1)/2} \mathcal{N} \left( \sigma^{(1-\alpha)/2} \sigma^{(1-\alpha)/2} \right) \mathcal{N}(\sigma)^{(\alpha-1)/2} \right]^{1/\alpha} = \mathcal{N}(\rho).
\]

(3.52)

Furthermore, for the same choices of \( \rho, \sigma, \) and \( \mathcal{N} \), (3.50) holds for \( \alpha \in (1, \infty) \).
Proof. Note that the quantities $\Delta_\alpha$ and $\tilde{\Delta}_\alpha$ defined in (1.15) and (1.16), respectively, can be expressed in terms of the $\alpha$-Rényi relative entropy and the $\alpha$-sandwiched Rényi relative entropy as follows \[SBW14\]:

\[
\Delta_\alpha (\rho, \sigma, N) = D_\alpha \left( \rho \left| \left| \sigma^{1-\alpha} N^{\frac{\alpha-1}{2}} N (\rho)^{\frac{1-\alpha}{2}} \sigma^{\frac{1-\alpha}{2}} \right| \right|^{\frac{1}{1-\alpha}} \right),
\]

(3.53)

and

\[
\tilde{\Delta}_\alpha (\rho, \sigma, N) = \tilde{D}_\alpha \left( \rho \left| \left| \sigma^{1-\alpha} N^{\frac{\alpha-1}{2}} N (\rho)^{\frac{1-\alpha}{2}} \sigma^{\frac{1-\alpha}{2}} \right| \right|^{\frac{1}{1-\alpha}} \right).
\]

(3.54)

The non-negativity and equality conditions of $\Delta_\alpha (\rho, \sigma, N)$ for $\alpha \in (0, 1)$, and that of $\tilde{\Delta}_\alpha (\rho, \sigma, N)$ for $\alpha \in (1/2, 1)$, then follow by applying Lemmas 1 and 8.

We next prove the non-negativity of $\Delta_\alpha (\rho, \sigma, N)$ for $\alpha \in (1, 2)$, $\rho \in S(H)_{++}$, $\sigma \in B(H)_{++}$, and $N$ a strict CPTP map. Using the definition (1.15) of $\Delta_\alpha (\rho, \sigma, N)$, cyclicity of trace, and the definition of the adjoint map, we can express $\Delta_\alpha (\rho, \sigma, N)$ as

\[
\Delta_\alpha (\rho, \sigma, N) = \frac{1}{\alpha - 1} \log \text{Tr} \left\{ N (\sigma)^{(\alpha-1)/2} N \left( \sigma^{(1-\alpha)/2} \rho^{\alpha} \sigma^{(1-\alpha)/2} \right) N (\sigma)^{(\alpha-1)/2} N (\rho)^{1-\alpha} \right\}
\]

(3.55)

\[
= D_\alpha \left( N (\sigma)^{(\alpha-1)/2} N \left( \sigma^{(1-\alpha)/2} \rho^{\alpha} \sigma^{(1-\alpha)/2} \right) N (\sigma)^{(\alpha-1)/2} \right)^{1/\alpha} \left| \left| N (\rho) \right| \right|
\]

(3.56)

where

\[
D_\alpha (P || Q) \equiv \frac{1}{\alpha - 1} \log \text{Tr} \left\{ P^{\alpha} Q^{1-\alpha} \right\}
\]

(3.57)

is the Rényi relative entropy between two positive definite operators $P$ and $Q$. It is known that $D_\alpha (P || Q) \geq 0$ if $P$ and $Q$ are such that

\[
\text{Tr} \left\{ P \right\} \geq \text{Tr} \left\{ Q \right\} = 1.
\]

(3.58)

This is because $D_\alpha (P || Q)$ is monotone under quantum operations for $\alpha \in (1, 2)$, and one such quantum operation is the trace operation:

\[
D_\alpha (P || Q) \geq D_\alpha \left( \text{Tr} \left\{ P \right\} || \text{Tr} \left\{ Q \right\} \right)
\]

(3.59)

\[
= \frac{1}{\alpha - 1} \log \left( \text{Tr} \left\{ P \right\}^{\alpha} \left| \text{Tr} \left\{ Q \right\} \right|^{1-\alpha} \right)
\]

(3.60)

\[
\geq 0.
\]

(3.61)

If $D_\alpha (P || Q) = 0$ for some $\alpha \in (1, 2)$ and $P$ and $Q$ such that $\text{Tr} \left\{ P \right\} \geq \text{Tr} \left\{ Q \right\}$, then it is known that $P = Q$ (one can deduce this, e.g., from the above and \[HMPB11\] Theorem 5.1). Hence, to prove the non-negativity of $\Delta_\alpha (\rho, \sigma, N)$ for $\alpha \in (1, 2)$, it suffices to prove that

\[
\text{Tr} \left\{ N (\sigma)^{(\alpha-1)/2} N \left( \sigma^{(1-\alpha)/2} \rho^{\alpha} \sigma^{(1-\alpha)/2} \right) N (\sigma)^{(\alpha-1)/2} \right\}^{1/\alpha} \geq \text{Tr} \left\{ N (\rho) \right\} = 1.
\]

(3.62)

Theorem 1.1 of \[Hia13\] establishes that the map

\[
(P, Q) \mapsto \text{Tr} \left\{ \left[ Q^{(\alpha-1)/2} P Q^{(\alpha-1)/2} \right]^{1/\alpha} \right\}
\]

(3.63)
is jointly concave for positive definite $P$ and $Q$ when $\alpha \in (1, 2)$. A straightforward argument allows to conclude its joint concavity for $\alpha \in (1, 2)$ and positive semidefinite $P$ and $Q$. Indeed, let $\varepsilon > 0$, 
$
\{P_x\}$ and $\{Q_x\}$ be sets of positive semidefinite operators, let $p_X(x)$ be a probability distribution, 
and let $P = \sum_x p_X(x) P_x$ and $\overline{Q} = \sum_x p_X(x) Q_x$. Consider that 
$$
\sum_x p_X(x) \text{Tr} \left\{ \left[ Q_x^{(\alpha-1)/2} P_x Q_x^{(\alpha-1)/2} \right]^{1/\alpha} \right\}
\leq \sum_x p_X(x) \text{Tr} \left\{ \left[ Q_x^{(\alpha-1)/2} (P_x + \varepsilon I) Q_x^{(\alpha-1)/2} \right]^{1/\alpha} \right\}
\leq \sum_x p_X(x) \text{Tr} \left\{ \left[ (P_x + \varepsilon I)^{1/2} Q_x^{\alpha-1} (P_x + \varepsilon I)^{1/2} \right]^{1/\alpha} \right\}
\leq \sum_x p_X(x) \text{Tr} \left\{ \left[ (P_x + \varepsilon I)^{1/2} (Q_x + \varepsilon I)^{\alpha-1} (P_x + \varepsilon I)^{1/2} \right]^{1/\alpha} \right\}
= \sum_x p_X(x) \text{Tr} \left\{ \left[ (Q_x + \varepsilon I)^{1/2} (P_x + \varepsilon I) (Q_x + \varepsilon I)^{1/2} \right]^{1/\alpha} \right\}
\leq \text{Tr} \left\{ \left[ (\overline{Q} + \varepsilon I)^{1/2} (P + \varepsilon I) (\overline{Q} + \varepsilon I)^{1/2} \right]^{1/\alpha} \right\}.
$$

The first inequality follows because $P_x \leq P_x + \varepsilon I$ and because $\text{Tr} \{ f(A) \} \leq \text{Tr} \{ f(B) \}$ for $A \leq B$ 
and $f(x) = x^{1/\alpha}$ a monotone non-decreasing function on $[0, \infty)$. The next inequality follows because 
$x^{\alpha-1}$ is operator monotone for $\alpha \in (1, 2)$ and for the same reason as above. The final inequality is 
a consequence of Theorem 1.1 of [Hal13]. By taking the limit $\varepsilon \searrow 0$, we can conclude that 
$$
\sum_x p_X(x) \text{Tr} \left\{ \left[ Q_x^{(\alpha-1)/2} P_x Q_x^{(\alpha-1)/2} \right]^{1/\alpha} \right\} \leq \text{Tr} \left\{ \left[ \overline{Q}^{(\alpha-1)/2} P \overline{Q}^{(\alpha-1)/2} \right]^{1/\alpha} \right\}.
$$

Now consider the following chain of equalities:

\begin{align*}
1 &= \text{Tr} \left\{ \left[ \sigma^{(\alpha-1)/2} \rho^{\alpha} \sigma^{(1-\alpha)/2} \sigma^{(\alpha-1)/2} \right]^{1/\alpha} \right\} \\
 &= \text{Tr} \left\{ \left[ \sigma^{(\alpha-1)/2} \rho^{\alpha} \sigma^{(1-\alpha)/2} \left( \sigma^{(1-\alpha)/2} \rho^{\alpha} \sigma^{(1-\alpha)/2} \otimes |0\rangle \langle 0|_E \right) \right]^{1/\alpha} \right\} \\
 &= \text{Tr} \left\{ \left[ \sigma^{(\alpha-1)/2} \otimes |0\rangle \langle 0|_E \right) \left( \sigma^{(1-\alpha)/2} \rho^{\alpha} \sigma^{(1-\alpha)/2} \otimes |0\rangle \langle 0|_E \right) \left( \sigma^{(\alpha-1)/2} \otimes |0\rangle \langle 0|_E \right) \right\}^{1/\alpha} \\
 &= \text{Tr} \left\{ \left[ U \sigma^{\frac{\alpha-1}{2}} \otimes |0\rangle \langle 0|_E \right) U^\dagger U \left( \sigma^{\frac{1-\alpha}{2}} \rho^{\alpha} \sigma^{\frac{1-\alpha}{2}} \otimes |0\rangle \langle 0|_E \right) U^\dagger \right\}^{1/\alpha} \\
 &= \text{Tr} \left\{ \left[ U \left( \sigma \otimes |0\rangle \langle 0|_E \right) U^\dagger \right]^{\frac{\alpha-1}{2}} \left[ U \left( \sigma \otimes |0\rangle \langle 0|_E \right) U^\dagger \right]^{\frac{\alpha-1}{2}} \right\}^{1/\alpha} \\
 &= \frac{1}{d_E^\alpha} \sum_i \text{Tr} \left\{ \left[ K_{i}^{\frac{\alpha-1}{2}} U_{E^i} U \left( \sigma^{1-\alpha} \rho^{\alpha} \sigma^{1-\alpha} \otimes |0\rangle \langle 0|_E \right) U^\dagger U_{E^i}^\dagger \left[ K_{i}^{\frac{\alpha-1}{2}} \right]^{1/2} \right]^{\frac{\alpha}{\alpha-1}} \right\} \\
\end{align*}

where 
$$
K_{i}^{\frac{\alpha-1}{2}} \equiv U_{E^i}^\dagger U \left( \sigma \otimes |0\rangle \langle 0|_E \right) U^\dagger U_{E^i}^\dagger,
$$

(3.66)
with \( \{ U_E^\dagger \} \) a set of Heisenberg-Weyl operators. Then by the concavity result (3.65) above for positive semidefinite \( P \) and \( Q \), it follows that the right-hand side of (3.66) is no larger than

\[
\text{Tr} \left\{ \left[ \mathcal{N}(\sigma) \otimes \pi_{E^\dagger} \right]^{(\alpha - 1)/2} \left[ \mathcal{N} \left( \sigma^{(1-\alpha)/2} \rho^\alpha \sigma^{(1-\alpha)/2} \right) \otimes \pi_{E^\dagger} \right] \left[ \mathcal{N}(\sigma) \otimes \pi_{E^\dagger} \right]^{(\alpha - 1)/2} \right\}^{1/\alpha}
\]

\[
= \text{Tr} \left\{ \mathcal{N}(\sigma)^{(\alpha - 1)/2} \mathcal{N} \left( \sigma^{(1-\alpha)/2} \rho^\alpha \sigma^{(1-\alpha)/2} \right) \mathcal{N}(\sigma)^{(\alpha - 1)/2} \right\}^{1/\alpha}. \tag{3.68}
\]

Thus we obtain the inequality

\[
\text{Tr} \left\{ \mathcal{N}(\sigma)^{(\alpha - 1)/2} \mathcal{N} \left( \sigma^{(1-\alpha)/2} \rho^\alpha \sigma^{(1-\alpha)/2} \right) \mathcal{N}(\sigma)^{(\alpha - 1)/2} \right\} \geq 1. \tag{3.69}
\]

This completes the proof of (3.48) for \( \alpha \in (1, 2) \). The equality condition in (3.52) follows from the representation in (3.50) and the equality condition stated after (3.61).

Next we prove that \( \Delta_\alpha(\rho, \sigma, \mathcal{N}) \geq 0 \) for \( \alpha \in (1, \infty) \). We start with the definition (1.16), which we repeat here for convenience:

\[
\Delta_\alpha(\rho, \sigma, \mathcal{N}) \equiv \frac{\alpha}{\alpha - 1} \log \left\| \rho^{1/2} \sigma^{1/2 - \alpha} \mathcal{N}^{\dagger} \left( \mathcal{N}(\sigma)^{\alpha-1/2} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{\alpha-1/2} \right) \sigma^{1/2 - \alpha} \rho^{1/2} \right\|_\alpha. \tag{3.70}
\]

From [MLDS+13, Lemma 12], it follows that the right-hand side of (3.70) can be written as

\[
\left\| \rho^{1/2} \sigma^{1/2 - \alpha} \mathcal{N}^{\dagger} \left( \mathcal{N}(\sigma)^{\alpha-1/2} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{\alpha-1/2} \right) \sigma^{1/2 - \alpha} \rho^{1/2} \right\|_\alpha
\]

\[
= \sup_{\tau > 0, \text{Tr} \{\tau\} \leq 1} \text{Tr} \left\{ \rho^{1/2} \sigma^{1/2 - \alpha} \mathcal{N}^{\dagger} \left( \mathcal{N}(\sigma)^{\alpha-1/2} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{\alpha-1/2} \right) \sigma^{1/2 - \alpha} \rho^{1/2} \tau^{1/\alpha} \right\}. \tag{3.71}
\]

Let us focus on the trace functional appearing on the right-hand side of the above equation:

\[
\Delta_\alpha(\rho, \sigma, \mathcal{N}; \tau)
\]

\[
\equiv \text{Tr} \left\{ \rho^{1/2} \sigma^{(1-\alpha)/2\alpha} \mathcal{N}^{\dagger} \left( \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \mathcal{N}(\rho)^{(1-\alpha)/\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \right) \sigma^{(1-\alpha)/2\alpha} \rho^{1/2} \tau^{(\alpha-1)/\alpha} \right\}
\]

\[
= \text{Tr} \left\{ \sigma^{(1-\alpha)/2\alpha} \rho^{1/2} \tau^{(\alpha-1)/\alpha} \sigma^{1/2 - \alpha} \mathcal{N}^{\dagger} \left( \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \mathcal{N}(\rho)^{(1-\alpha)/\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \right) \right\}
\]

\[
= \text{Tr} \left\{ \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \mathcal{N}(\sigma)^{(1-\alpha)/2\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \right\} \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \mathcal{N}(\sigma)^{(1-\alpha)/\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \right\}
\]

\[
= \text{Tr} \left\{ \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \mathcal{N}(\sigma)^{(1-\alpha)/2\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \mathcal{N}(\sigma)^{(1-\alpha)/\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2\alpha} \right\}. \tag{3.72}
\]

By making the substitution \( \alpha' \equiv (2\alpha - 1) / \alpha \), so that \( (1 - \alpha) / \alpha = 1 - \alpha' \) and thus \( \alpha' \in (1, 2) \) when \( \alpha \in (1, \infty) \), we see that the last line above is equal to

\[
\text{Tr} \left\{ \mathcal{N}(\sigma)^{(\alpha'-1)/2} \mathcal{N}(\sigma^{(1-\alpha')/2}) \rho^{1/2} \sigma^{1/2} \mathcal{N}(\sigma)^{(\alpha'-1)/2} \mathcal{N}(\rho)^{(1-\alpha')/2} \right\}. \tag{3.73}
\]

Observe that this is similar to (3.55). Hence we can write \( \Delta_\alpha(\rho, \sigma, \mathcal{N}; \tau) \) as

\[
\frac{1}{\alpha' - 1} \log \text{Tr} \left\{ \mathcal{N}(\sigma)^{(\alpha'-1)/2} \mathcal{N}(\sigma^{(1-\alpha')/2}) \rho^{1/2} \sigma^{1/2} \mathcal{N}(\sigma)^{(\alpha'-1)/2} \mathcal{N}(\rho)^{(1-\alpha')/2} \right\}
\]

\[
= D_{\alpha'} \left[ \mathcal{N}(\sigma)^{(\alpha'-1)/2} \mathcal{N}(\sigma^{(1-\alpha')/2}) \rho^{1/2} \sigma^{1/2} \mathcal{N}(\sigma)^{(\alpha'-1)/2} \mathcal{N}(\rho)^{(1-\alpha')/2} \right]^{1/\alpha'} \left\| \mathcal{N}(\rho) \right\|. \tag{3.74}
\]
This implies that

\[
\Delta_\alpha (\rho, \sigma, N)
= \sup_{\tau \geq 0} \left\{ \left[ N(\sigma)^{\alpha' - 1/2} N\left( (1 - \sigma) / \rho^{1/2} (\alpha' - 1/2)(1 - \sigma)^{1/2} \right) N(\sigma)^{\alpha' - 1/2} \right]^{1/\alpha'} \right\}
\geq 0,
\]

where the first inequality follows by setting \( \tau = \rho \) and the last inequality follows because we have already shown that \( \Delta_{\alpha'} (\rho, \sigma, N) \geq 0 \) for \( \alpha' \in (1, 2) \). The above inequality also demonstrates that

\[
\Delta_\alpha (\rho, \sigma, N) \geq \Delta_{\alpha'} (\rho, \sigma, N).
\]
Hence,
\[ I_\alpha(A; B|C)_\rho = \frac{1}{\alpha - 1} \log \sum_j q(j) = 0. \] (4.2)

Now we prove the second statement. From Lemma 10 we have that
\[ \text{Tr} \left\{ \left( \rho_{AC}^{(1-\alpha)/2} \rho_{BC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right)^{1/(1-\alpha)} \right\} \leq 1 \] (4.3)
for \( \alpha \in (0, 1) \cup (1, 2) \). Suppose that
\[ I_\alpha(A; B|C)_\rho = 0, \] (4.4)
for some \( \alpha \in (0, 1) \). Recall from [BSW15] that
\[ I_\alpha(A; B|C)_\rho = D_\alpha \left( \rho_{ABC} \left| \left| \left( \rho_{AC}^{(1-\alpha)/2} \rho_{BC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right)^{1/(1-\alpha)} \right) \right. \right), \] (4.5)
where \( D_\alpha \) is the \( \alpha \)-Rényi relative entropy. Applying Lemmas 1 and 10, we find that
\[ \rho_{ABC} = \left( \rho_{AC}^{(1-\alpha)/2} \rho_{BC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right)^{1/(1-\alpha)} \]
\[ \Leftrightarrow \rho_{ABC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} = \rho_{AC}^{(1-\alpha)/2} \rho_{BC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \]
\[ \Leftrightarrow \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} = \rho_{AC}^{(1-\alpha)/2} \rho_{BC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \] (4.6)
We can then multiply both sides of the last line by \( \rho_{AC} \) and take the trace of the last line to find that
\[ \text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right\} = \text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right\} \]
\[ \Leftrightarrow \text{Tr} \left\{ \rho_{ABC} \rho_{AC}^{(1-\alpha)/2} \right\} = \text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right\} \]
\[ \Leftrightarrow \text{Tr} \left\{ \rho_{ABC} \rho_{AC}^{(1-\alpha)/2} \right\} = \text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right\}. \] (4.7)
Recognizing the last line as an equality of \( f \)-divergences (see Remark 3), it is then equivalent to (3.15) of Lemma 4, which by (3.16) of the lemma implies that
\[ \log \rho_{ABC} = \log \rho_{AC} + \log \rho_{BC} - \log \rho_{C}. \] (4.8)
By Lemma 5 we then conclude that \( \rho_{ABC} \) is a short quantum Markov chain \( A - C - B \).

To get the second statement of the theorem for some \( \alpha \in (1, 2) \), by the same reasoning as above, we can obtain (4.6), i.e.,
\[ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} = \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2}. \] (4.9)
Taking the matrix inverse of both sides then gives
\[ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} = \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2}. \] (4.10)
Multiplying both sides by \( \rho_{AC} \) and taking the trace gives
\[ \text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right\} = \text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right\}, \] (4.11)
which simplifies to
\[ \text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right\} = \text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right\}. \] (4.12)
The last line is once again an equality of \( f \)-divergences, which by Lemmas 4 and 5 implies that \( \rho_{ABC} \) is a short quantum Markov chain \( A - C - B \).
5 Sandwiched Rényi conditional mutual information and short quantum Markov chains

Theorem 14 Let $\rho_{ABC} \in \mathcal{S}(\mathcal{H}_{ABC})_{++}$. Then $\overline{I}_\alpha(A;B|C)_\rho = 0$ for all $\alpha \in (0,1) \cup (1,\infty)$ if $\rho_{ABC}$ is a short quantum Markov chain $A-C-B$. Furthermore, $\rho_{ABC}$ is a short quantum Markov chain $A-C-B$ if $\overline{I}_\alpha(A;B|C)_\rho = 0$ for some $\alpha \in (1/2,1) \cup (1,\infty)$.

The proof of the above theorem is analogous to that of Theorem 13. For the sake of completeness, we include it in Appendix A.

6 Sufficiency of quantum channels and Rényi generalizations of relative entropy differences

Theorem 15 Let $\rho \in \mathcal{S}(\mathcal{H})_{++}$, $\sigma \in \mathcal{B}(\mathcal{H})_{++}$, and let $\mathcal{N}$ be a strict CPTP map. Then $\Delta_\alpha(\rho,\sigma,\mathcal{N}) = 0$ and $\overline{\Delta}_\alpha(\rho,\sigma,\mathcal{N}) = 0$ for all $\alpha \in (0,1) \cup (1,\infty)$ if $\mathcal{N}$ is sufficient for $\rho$ and $\sigma$. Furthermore, $\mathcal{N}$ is sufficient for $\rho$ and $\sigma$ if $\Delta_\alpha(\rho,\sigma,\mathcal{N}) = 0$ for some $\alpha \in (0,1) \cup (1,2)$ or if $\overline{\Delta}_\alpha(\rho,\sigma,\mathcal{N})$ for some $\alpha \in (1/2,1) \cup (1,\infty)$.

Proof. Suppose that $\mathcal{N}$ is sufficient for $\rho$ and $\sigma$. Then according to Theorem 3 this implies that the decompositions in (3.19)-(3.20) hold. It suffices to consider the contribution arising from a given block (as in our previous proofs), so we focus on this case. We consider the quantity $\Delta_\alpha(\rho,\sigma,\mathcal{N})$ first. Then for a given block, evaluating the formula

$$\rho^{\alpha}\sigma^{(1-\alpha)/2}\mathcal{N}^{\dagger}\left(\mathcal{N}(\sigma)^{(\alpha-1)/2}\mathcal{N}(\rho)^{1-\alpha}\mathcal{N}(\sigma)^{(\alpha-1)/2}\right)^{\sigma^{(1-\alpha)/2}} (6.1)$$

gives

$$(pp_L \otimes \tau_R)^{\alpha}(q\sigma_L \otimes \tau_R)^{(1-\alpha)/2} \times$$

$$\left[\mathcal{U}^{\dagger} \otimes (\mathcal{N}^R)^{\dagger}\right] \left(\left(q\mathcal{U}(\sigma_L) \otimes \mathcal{N}^R(\tau_R)\right)^{\alpha-1/2} \left(p\mathcal{U}(\rho_L) \otimes \mathcal{N}^R(\tau_R)\right)^{(1-\alpha)/2}\right) \times$$

$$(q\sigma_L \otimes \tau_R)^{(1-\alpha)/2}$$

$$\frac{p}{\left(\rho_L \otimes \tau_R\right)^{\alpha}(\sigma_L \otimes \tau_R)^{(1-\alpha)/2}} \times$$

$$\left[\mathcal{U}^{\dagger} \otimes (\mathcal{N}^R)^{\dagger}\right] \left(\left(U(\sigma_L) \otimes \mathcal{N}^R(\tau_R)\right)^{(\alpha-1)/2} \left(U(\rho_L) \otimes \mathcal{N}^R(\tau_R)\right)^{1-\alpha}(\sigma_L \otimes \mathcal{N}^R(\tau_R))^{(\alpha-1)/2}\right) \times$$

$$(\sigma_L \otimes \tau_R)^{(1-\alpha)/2} \ (6.2)$$

$$= p\left(\rho_L \otimes \tau_R\right)^{\alpha}(\sigma_L \otimes \tau_R)^{(1-\alpha)/2} \times$$

$$\left[\mathcal{U}^{\dagger} \otimes (\mathcal{N}^R)^{\dagger}\right] \left(\left(U(\sigma_L)^{\alpha-1/2} \mathcal{U}(\rho_L) \otimes \mathcal{N}^R(\tau_R)^{\alpha-1/2}\right)^{\alpha-1/2} \right) \times$$

$$(\sigma_L \otimes \tau_R)^{(1-\alpha)/2} \ (6.3)$$

18
Continuing, the last line is equal to
\[
p \left( (\rho_L)^{\alpha} (\sigma_L)^{(1-\alpha)/2} \otimes \tau_R^{(\alpha+1)/2} \right) \times \\
\left[ \mathcal{U}^d \otimes (\mathcal{N}^R)^{1/4} \right] \left( \mathcal{U} \left[ (\sigma_L)^{(\alpha+1)/2} (\rho_L)^{1-\alpha} (\sigma_L)^{(\alpha-1)/2} \right] \otimes \mathcal{I}_R \right) (\sigma_L \otimes \tau_R)^{(1-\alpha)/2} \\
= p \left( (\rho_L)^{\alpha} (\sigma_L)^{(1-\alpha)/2} \otimes \tau_R^{(\alpha+1)/2} \right) \times \\
\left[ (\sigma_L)^{(\alpha-1)/2} (\rho_L)^{1-\alpha} (\sigma_L)^{(\alpha-1)/2} \right] \otimes \mathcal{I}_R \right) (\sigma_L \otimes \tau_R)^{(1-\alpha)/2} \\
= p \left( (\rho_L)^{\alpha} (\sigma_L)^{(1-\alpha)/2} (\rho_L)^{1-\alpha} (\sigma_L)^{(\alpha-1)/2} (\sigma_L)^{(1-\alpha)/2} \otimes \tau_R^{(\alpha+1)/2} \tau_R^{(1-\alpha)/2} \right) \\
= p \left( (\rho_L)^{\alpha} (\rho_L)^{1-\alpha} \otimes \tau_R \right) \\
= \rho \rho_L \otimes \tau_R. \tag{6.4}
\]

Taking the trace we find that each block has trace equal to \( p(j) \), so that
\[
\Delta_\alpha (\rho, \sigma, \mathcal{N}) = \frac{1}{\alpha - 1} \log \sum_j p(j) = 0. \tag{6.5}
\]

So this proves that \( \Delta_\alpha (\rho, \sigma, \mathcal{N}) = 0 \) for all \( \alpha \in (0, 1) \cup (1, \infty) \) if \( \mathcal{N} \) is sufficient for \( \rho \) and \( \sigma \). The statement for \( \tilde{\Delta}_\alpha (\rho, \sigma, \mathcal{N}) \) follows by combining the above development with that in the first part of the proof of Theorem \( \text{14} \).

Now suppose that \( \Delta_\alpha (\rho, \sigma, \mathcal{N}) = 0 \) for some \( \alpha \in (0, 1) \). From Proposition \( \text{12} \) we have that
\[
\rho = \left[ \sigma^{(1-\alpha)/2} \mathcal{N}^\dagger \left( \mathcal{N}(\sigma)^{(\alpha-1)/2} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2} \right) \sigma^{(1-\alpha)/2} \right]^{1/(1-\alpha)}, \tag{6.6}
\]
which, for positive definite \( \sigma \), is equal to
\[
\sigma^{(\alpha-1)/2} \rho^{1-\alpha} \sigma^{(\alpha-1)/2} = \mathcal{N}^\dagger \left( \mathcal{N}(\sigma)^{(\alpha-1)/2} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2} \right). \tag{6.7}
\]

Multiply both sides by \( \sigma \) and take the trace to get
\[
\text{Tr} \left\{ \sigma \sigma^{(\alpha-1)/2} \rho^{1-\alpha} \sigma^{(\alpha-1)/2} \right\} = \text{Tr} \left\{ \sigma \mathcal{N}^\dagger \left( \mathcal{N}(\sigma)^{(\alpha-1)/2} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2} \right) \right\} \\
\iff \text{Tr} \left\{ \rho^{1-\alpha} \sigma^{\alpha} \right\} = \text{Tr} \left\{ \mathcal{N}(\sigma)^{(\alpha-1)/2} \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{(\alpha-1)/2} \right\} \\
\iff \text{Tr} \left\{ \rho^{1-\alpha} \sigma^{\alpha} \right\} = \text{Tr} \left\{ \mathcal{N}(\rho)^{1-\alpha} \mathcal{N}(\sigma)^{\alpha} \right\} \tag{6.8}
\]

We can then conclude that the conditions in \( \text{3.19} \)–\( \text{3.20} \) hold by the same sequence of reasoning as in previous proofs, so that \( \mathcal{N} \) is sufficient for \( \rho \) and \( \sigma \).

If for some \( \alpha \in (1, 2) \) we have \( \Delta_\alpha (\rho, \sigma, \mathcal{N}) = 0 \), then we know that
\[
\left[ \mathcal{N}(\sigma)^{(\alpha-1)/2} \mathcal{N} \left( \sigma^{(\alpha-1)/2} \rho^{1-\alpha} \sigma^{(\alpha-1)/2} \right) \mathcal{N}(\sigma)^{(\alpha-1)/2} \right]^{1/\alpha} = \mathcal{N}(\rho) \tag{6.9}
\]
\[
\iff \mathcal{N} \left( \sigma^{(\alpha-1)/2} \rho^{1-\alpha} \sigma^{(\alpha-1)/2} \right) = \mathcal{N}(\rho)^{(1-\alpha)/2} \mathcal{N}(\rho)^{\alpha} \mathcal{N}(\sigma)^{(1-\alpha)/2} \tag{6.10}
\]
This implies that
\[
\text{Tr} \left\{ \mathcal{N} \left( \sigma^{(1-\alpha)/2} \rho^\alpha \sigma^{(1-\alpha)/2} \right) \right\} = \text{Tr} \left\{ \mathcal{N} (\rho)^\alpha \mathcal{N} (\sigma)^{1-\alpha} \right\}
\]
(6.11)
which in turn implies that
\[
\text{Tr} \left\{ \rho^\alpha \sigma^{1-\alpha} \right\} = \text{Tr} \left\{ \mathcal{N} (\rho)^\alpha \mathcal{N} (\sigma)^{1-\alpha} \right\},
\]
(6.12)
because \( N \) is trace preserving. This is then an equality of \( f \)-divergences, which we can use to conclude the sufficiency property as in previous proofs.

If \( \bar{\Delta}_\alpha (\rho, \sigma, N) = 0 \), for some \( \alpha \in (1, \infty) \), then from Proposition [12] and (3.80) we have that \( \Delta_{\alpha'} (\rho, \sigma, N) = 0 \) for some \( \alpha' \in (1, 2) \). Then we know from the above analysis that \( N \) is sufficient for \( \rho \) and \( \sigma \). ■

7 Quantum Markov chains, sufficiency of quantum channels, and min- and max-information measures

7.1 Min- and max- generalizations of conditional mutual information

Definition 16 ([BSW15]) For a tripartite state \( \rho_{ABC} \in S (\mathcal{H}_{ABC})_{++} \) the max-conditional mutual information and min-conditional mutual information are defined respectively as follows:

\[
I_{\text{max}} (A; B | C) \equiv D_{\text{max}} \left( \rho_{ABC} \left\| \rho_{AC}^{1/2} \rho_{BC}^{1/2} \right\| \rho_{AC}^{-1/2} \rho_{BC}^{-1/2} \right)
\]
(7.1)
\[
= \inf \left\{ \lambda : \rho_{ABC} \leq \exp (\lambda) \rho_{AC}^{1/2} \rho_{BC}^{1/2} \rho_{AC}^{-1/2} \rho_{BC}^{-1/2} \right\}
\]
(7.2)
\[
= 2 \log \left\| \rho_{ABC} \rho_{AC}^{1/2} \rho_{BC}^{1/2} \rho_{AC}^{-1/2} \rho_{BC}^{-1/2} \right\|_{\infty},
\]
(7.3)
\[
I_{\text{min}} (A; B | C) \equiv D_{\text{min}} \left( \rho_{ABC} \left\| \rho_{AC}^{1/2} \rho_{BC}^{1/2} \right\| \rho_{AC}^{-1/2} \rho_{BC}^{-1/2} \right)
\]
(7.4)
\[
= - \log F \left( \rho_{ABC}, \rho_{AC}^{1/2} \rho_{BC}^{1/2} \rho_{AC}^{-1/2} \rho_{BC}^{-1/2} \right)
\]
(7.5)
\[
= - 2 \log \left\| \rho_{ABC} \rho_{AC}^{1/2} \rho_{BC}^{1/2} \rho_{AC}^{-1/2} \rho_{BC}^{-1/2} \right\|_{1}.
\]
(7.6)

As observed in [BSW15], we have that
\[
I_{\text{max}} (A; B | C)_\rho, I_{\text{min}} (A; B | C)_\rho \geq 0,
\]
(7.7)
due to Lemma 1 and the fact that \( \text{Tr} \left\{ \rho_{AC}^{1/2} \rho_{BC}^{1/2} \rho_{AC}^{-1/2} \rho_{BC}^{-1/2} \right\} = 1 \). The following theorem gives the equality conditions for (7.7):

Theorem 17 Let \( \rho_{ABC} \in S (\mathcal{H}_{ABC})_{++} \). Then each of the following identities hold if and only if \( \rho_{ABC} \) is a short quantum Markov chain \( A - C - B \):

\[
I_{\text{max}} (A; B | C) = 0, \quad I_{\text{min}} (A; B | C) = 0.
\]
(7.8)
Note that the right hand side of the above equation is a state obtained by the action of the Petz\ operator $I$ is a density operator. Then from (3.13) of Lemma 1 it follows that

$$Hence, that of the expression for $I$ of the operators $A$ and the relation
to $\|A\|_\infty$ arising from the $j$th block: For $I_{max}(A;B|C)$ this gives the following contribution (where we again suppress the index $j$ for simplicity) to the infinity norm appearing in (7.3):

$$\left\| (q\rho_{AC} \otimes \rho_{R}B)^{1/2} (q\rho_{AC} \otimes \rho_{C})^{-1/2} (q\rho_{C} \otimes \rho_{R})^{1/2} (q\rho_{C} \otimes \rho_{C}B)^{-1/2} \right\|_\infty$$

$$= \left\| \left( \rho_{AC}^{1/2} \otimes \rho_{C}^{-1/2} \right) \left( \rho_{C}^{-1/2} \otimes \rho_{C}^{1/2} \right) \left( \rho_{C}^{1/2} \otimes \rho_{C}^{-1/2} \right) \left( \rho_{C}^{-1/2} \otimes \rho_{C}^{1/2} \right) \right\|_\infty$$

$$= q \left\| \left( \rho_{AC}^{1/2} \otimes \rho_{C}^{-1/2} \right) \left( \rho_{C}^{-1/2} \otimes \rho_{C}^{1/2} \right) \left( \rho_{C}^{1/2} \otimes \rho_{C}^{-1/2} \right) \left( \rho_{C}^{-1/2} \otimes \rho_{C}^{1/2} \right) \right\|_1$$

$$= q \left\| \rho_{AC} \otimes \rho_{C}B \right\|_1$$

This, along with the relation $\|a_1A_1 \oplus a_2A_2\|_2 = a_1 \|A_1\|_1 + a_2 \|A_2\|_1$ for scalars $a_1$ and $a_2$ and operators $A_1$ and $A_2$, yields

$$I_{max}(A;B|C) = -2 \log \left( \sum_j q(j) \right) = 0.$$

To prove the converse, we first start with the assumption that $I_{max}(A;B|C) = 0$ and make use of the expression for $I_{max}(A;B|C)$ in terms of the max-relative entropy, as given by (7.1). Note that

$$\text{Tr} \left( \rho_{AC}^{-1/2} \rho_{BC}^{-1/2} \rho_{C}^{-1/2} \rho_{AC}^{-1/2} \right) = 1.$$  

Hence,

$$\tau_{ABC} \equiv \rho_{AC}^{-1/2} \rho_{BC}^{-1/2} \rho_{C}^{-1/2} \rho_{AC}^{-1/2}$$

is a density operator. Then from (3.13) of Lemma 1 it follows that $I_{max}(A;B|C) = 0$ if and only if

$$\rho_{ABC} = \rho_{AC}^{-1/2} \rho_{BC}^{-1/2} \rho_{C}^{-1/2} \rho_{AC}^{-1/2}.$$
mutual information $I(A;B|C)$ of $\rho_{ABC}$ to be equal to zero. This in turn (by [HJPW04]) implies that $\rho_{ABC}$ is a short quantum Markov chain $A - C - B$.

We next assume that $I_{\min}(A;B|C)_\rho = 0$. Note that $I_{\min}(A;B|C)_\rho$ is expressible as the min-relative entropy (as given in (7.4)) of two density operators $\rho_{ABC}$ and $\tau_{ABC}$ (defined in (7.20)). Then (3.12) of Lemma 1 implies that the identity (7.21) holds. From the results of [HJPW04], this establishes $\rho_{ABC}$ to be a short quantum Markov chain $A - C - B$.

7.2 Min- and max- generalizations of a relative entropy difference

We consider the following generalizations of a relative entropy difference, motivated by the developments in [BSW15, SBW14].

**Definition 18** Let $\rho \in S(\mathcal{H})_{++}$, $\sigma \in B(\mathcal{H})_{++}$, and let $\mathcal{N}$ be a strict CPTP map. Then,

$$\Delta_{\min} (\rho, \sigma, \mathcal{N}) \equiv D_{\min} (\rho || R_{\sigma, \mathcal{N}}(\mathcal{N}(\rho))) ,$$

(7.22)

where $R_{\sigma, \mathcal{N}}$ is the Petz recovery channel defined in (1.10) and

$$\Delta_{\max} (\rho, \sigma, \mathcal{N}) \equiv D_{\max} (\rho || R_{\sigma, \mathcal{N}}(\mathcal{N}(\rho))) .$$

(7.23)

**Theorem 19** For $\rho \in S(\mathcal{H})_{++}$, $\sigma \in B(\mathcal{H})_{++}$, and a strict CPTP map $\mathcal{N}$,

$$\Delta_{\min} (\rho, \sigma, \mathcal{N}) \geq 0, \quad \Delta_{\max} (\rho, \sigma, \mathcal{N}) \geq 0,$$

(7.24)

with equality holding if and only if $\mathcal{N}$ is sufficient for $\rho$ and $\sigma$.

**Proof.** The non-negativity conditions follow from the fact that $R_{\sigma, \mathcal{N}}(\mathcal{N}(\rho))$ is a density operator (since $R_{\sigma, \mathcal{N}}$ is a CPTP map) and Lemma 1. The equality conditions also follow from Lemma 1 and the fact that if $\mathcal{N}$ is sufficient for $\rho$ and $\sigma$ then $R_{\sigma, \mathcal{N}}(\mathcal{N}(\rho)) = \rho$. That $\mathcal{N}$ is sufficient for $\rho$ and $\sigma$ implies the equality condition follows from a proof similar to that in Theorem 17, but employing the direct-sum structure from Theorem 6. Furthermore, if $\Delta_{\max} (\rho, \sigma, \mathcal{N}) = 0$ or $\Delta_{\min} (\rho, \sigma, \mathcal{N}) = 0$, then Lemma 1 implies that

$$R_{\sigma, \mathcal{N}}(\mathcal{N}(\rho)) = \rho,$$

(7.25)

which allows us to conclude that $\mathcal{N}$ is sufficient for $\rho$ and $\sigma$.

8 Conclusion and open questions

We have shown that the $\alpha$-Rényi conditional mutual informations $I_{\alpha}(A;B|C)_\rho$ and $\tilde{I}_{\alpha}(A;B|C)_\rho$ from [BSW15] are equal to zero if and only if $\rho_{ABC}$ is a short quantum Markov chain $A - C - B$. We have also shown that the $\alpha$-Rényi quantities $\Delta_{\alpha}(\rho, \sigma, \mathcal{N})$ and $\tilde{\Delta}_{\alpha}(\rho, \sigma, \mathcal{N})$ from [SBW14] are non-negative and equal to zero if and only if $\mathcal{N}$ is sufficient for $\rho$ and $\sigma$. Moreover, we have solved some open questions from [Rus02, Zha14].

There are some interesting open questions to consider going forward from here. We would like to know the equality conditions for monotonicity of the sandwiched Rényi relative entropies, i.e., for which triples $(\rho, \sigma, \mathcal{N})$ is it true that

$$\tilde{D}_{\alpha}(\rho || \sigma) = \tilde{D}_{\alpha}(\mathcal{N}(\rho) || \mathcal{N}(\sigma)) ?$$

(8.1)
Hence, direct-sum decomposition (1.7). This gives the following (where we suppress the \( j \) in [FR14, BLW14] might suggest.

\[ \rho \]

\equal{} zero? Presumably the former has to do with \( \rho, \sigma, \) is there a characterization of triples \((\rho, \sigma, N)\) for which \( \Delta_\alpha (\rho, \sigma, N) \) and \( \overline{\Delta}_\alpha (\rho, \sigma, N) \) are nearly equal to zero? Presumably the former has to do with \( \rho_{ABC} \) being close to the “Petz recovered” \( \rho_{BC} \) and the latter has to do with \( \rho \) being close to the “Petz recovered” \( \rho \), as recent developments in [FR14, BLW14] might suggest.

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A Proof of Theorem 14

Proof. To prove the first statement of Theorem 14, suppose that \( \rho_{ABC} \) is a short quantum Markov chain \( A - C - B \) so that (1.6)-(1.7) hold. We can then directly evaluate the formula in (1.14). As in the previous proofs, it suffices to evaluate the contribution arising from each block \( j \) in the direct-sum decomposition (1.7). This gives the following (where we suppress the \( j \) index, again for simplicity)

\[
\left\| (q\rho_{AC_L} \otimes \rho_{CR_B})^{1/2} \left( (q\rho_{AC_L} \otimes \rho_{CR})^{(1-\alpha)/2\alpha} (q\rho_{CL} \otimes \rho_{CR})^{(\alpha-1)/2\alpha} (q\rho_{CL} \otimes \rho_{CR})^{(1-\alpha)/2\alpha} \right)^{2\alpha} \right\|_{2\alpha}^2 \\
= q \left\| \left( \rho_{AC_L} \otimes \rho_{CR_B} \right)^{1/2} \left( \rho_{AC_L}^{(1-\alpha)/2\alpha} \otimes \rho_{CR_B}^{(\alpha-1)/2\alpha} \right) \left( \rho_{CL}^{(\alpha-1)/2\alpha} \otimes \rho_{CR}^{(\alpha-1)/2\alpha} \right) \left( \rho_{CL}^{(1-\alpha)/2\alpha} \otimes \rho_{CR}^{(1-\alpha)/2\alpha} \right) \right\|_{2\alpha}^{2\alpha} \\
= q \left\| \left( \rho_{AC_L}^{1/2\alpha} \rho_{AC_L}^{(1-\alpha)/2\alpha} \rho_{CL}^{(\alpha-1)/2\alpha} \rho_{CL}^{(1-\alpha)/2\alpha} \right) \left( \rho_{CR_B}^{1/2\alpha} \rho_{CR_B}^{(1-\alpha)/2\alpha} \rho_{CR}^{(\alpha-1)/2\alpha} \rho_{CR}^{(1-\alpha)/2\alpha} \right) \right\|_{2\alpha}^{2\alpha} \\
= q \left\| \rho_{AC_L}^{1/2\alpha} \rho_{CR_B}^{1/2\alpha} \right\|_{2\alpha}^{2\alpha} \\
= q. \tag{A.1}
\]

Hence,

\[
\overline{I}_\alpha (A; B | C) = \frac{1}{\alpha - 1} \log \left( \sum_j q(j) \right) = 0, \tag{A.3}
\]
where we used that $\|A \oplus B\|_p^p = \|A\|_p^p + \|B\|_p^p$.

Now we prove the second statement. From Corollary 10 we have that
\[
\text{Tr} \left\{ \left( \rho_{AC}^{(1-\alpha)/2} \rho_C^{(1-\alpha)/2} \rho_{BC}^{(1-\alpha)/2} \rho_C^{(1-\alpha)/2} \right) \right\} \leq 1 \tag{A.4}
\]
for $\alpha \in (1/2, 1) \cup (1, \infty)$. So, suppose that
\[
\tilde{I}_\alpha(A; B|C)_\rho = 0,
\]
for some $\alpha \in (1/2, 1)$. Recall from [BSW15] that
\[
\tilde{I}_\alpha(A; B|C)_\rho = D_\alpha \left( \rho_{ABC} \left| \rho_{AC}^A \rho_C^B \right. \right), \tag{A.6}
\]
where $\tilde{D}_\alpha$ is the sandwiched Rényi relative entropy [MLDS13, WWY14]. Applying Lemmas 1 and 10 we find that
\[
\rho_{ABC} = \left( \rho_{AC}^{(1-\alpha)/2} \rho_C^{(1-\alpha)/2} \rho_{BC}^{(1-\alpha)/2} \rho_C^{(1-\alpha)/2} \right)^{\alpha/(1-\alpha)}.
\]
We can then multiply both sides of the last line by $\rho_{AC}$ and then take the trace to find that
\[
\text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{ABC} \rho_{AC}^{(1-\alpha)/2} \right\} = \text{Tr} \left\{ \rho_C^{(1-\alpha)/2} \rho_{BC} \rho_C^{(1-\alpha)/2} \rho_C^{(1-\alpha)/2} \right\}
\]
\[
\Leftrightarrow \text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right\} = \text{Tr} \left\{ \rho_C^{(1-\alpha)/2} \rho_{BC} \rho_C^{(1-\alpha)/2} \rho_C^{(1-\alpha)/2} \right\}
\]
\[
\Leftrightarrow \text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right\} = \text{Tr} \left\{ \rho_C^{(1-\alpha)/2} \rho_{BC} \rho_C^{(1-\alpha)/2} \rho_C^{(1-\alpha)/2} \right\}. \tag{A.8}
\]
For $\alpha \in (1/2, 1)$, this statement is then equivalent to (3.15) of Lemma 4 which by (3.16) of the lemma implies that
\[
\log \rho_{ABC} = \log \rho_{AC} + \log \rho_{BC} - \log \rho_C. \tag{A.9}
\]
By Lemma 5 we can then infer that $\rho_{ABC}$ is a short quantum Markov chain $A - C - B$. To get the second statement for some $\alpha \in (1, \infty)$, by the same reasoning as above, we can conclude (A.7), i.e.,
\[
\rho_{AC}^{(1-\alpha)/2} \rho_{ABC} \rho_{AC}^{(1-\alpha)/2} = \rho_C^{(1-\alpha)/2} \rho_{BC} \rho_C^{(1-\alpha)/2} \rho_C^{(1-\alpha)/2}.
\]
Taking the matrix inverse of both sides then gives
\[
\rho_{AC}^{(1-\alpha)/2} \rho_{ABC} \rho_{AC}^{(1-\alpha)/2} = \rho_C^{(1-\alpha)/2} \rho_{BC} \rho_C^{(1-\alpha)/2} \rho_C^{(1-\alpha)/2}.
\]
Multiplying both sides by $\rho_{AC}$ and taking the trace, we get
\[
\text{Tr} \left\{ \rho_{AC}^{(1-\alpha)/2} \rho_{ABC} \rho_{AC}^{(1-\alpha)/2} \rho_C^{(1-\alpha)/2} \rho_{AC}^{(1-\alpha)/2} \right\} = \text{Tr} \left\{ \rho_C^{(1-\alpha)/2} \rho_{BC} \rho_C^{(1-\alpha)/2} \rho_C^{(1-\alpha)/2} \right\}, \tag{A.12}
\]
which simplifies to
\[
\text{Tr} \left\{ \rho_{ABC} \rho_{AC}^{1/\alpha} \right\} = \text{Tr} \left\{ \rho_{BC} \rho_C^{1/\alpha} \right\}. \tag{A.13}
\]
The last line is once again an equality of $f$-divergences, which by Lemmas 4 and 5 implies that $\rho_{ABC}$ is a short quantum Markov chain $A - C - B$. ■

24
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