Abstract—We show that many delay-based reservoir computers considered in the literature can be characterized by a universal master memory function (MMF). Once computed for two independent parameters, this function provides linear memory capacity for any delay-based single-variable reservoir with small inputs. Moreover, we propose an analytical description of the MMF that enables its efficient and fast computation. Our approach can be applied not only to single-variable delay-based reservoirs governed by known dynamical rules, such as the Mackey–Glass or Stuart–Landau-like systems, but also to reservoirs whose dynamical model is not available.

Index Terms—Machine learning, nonlinear dynamics, reservoir computing.

I. INTRODUCTION

RESERVOIR computing is a neuromorphic-inspired machine learning paradigm, which enables high-speed training of recurrent neural networks and is capable of solving highly complex time-dependent tasks. First proposed by Jaeger [1] and inspired by the human brain [2], it utilizes the inherent computational capabilities of dynamical systems. Very recently, the universal approximation property has also been shown for a wide range of reservoir computers, which solidifies the concept as a broadly applicable scheme [3]. Bollt et al. [4], Gauthier et al. [5], and Jauring and Lüdge, [6] pointed out a connection between reservoir computers and vector autoregressive (VAR) and nonlinear VAR machines, which may be one of the reasons behind the surprising efficiency of reservoir computers for time-dependent tasks. Realizations range from electrical and optoelectrical [7], [8], [9] up to optical setups [10], [11], [12], [13], [14] and have shown the relevance of reservoir computing to practical applications like human action recognition [15]. Additionally, analytical and numerical analyses [16], [17], [18], [19] help in building an understanding of its working principles and improve its performance while pinpointing to easily implementable realizations [20], [21]. Utilizing reservoir computers as a handy supporting tool for predicting difficult spatiotemporal patterns has shown promising results in increasing the accuracy for predictions of chaotic dynamics like weather forecasting [22], [23]. Many groups endeavor to optimize delay-based reservoir computing, e.g., through the utilization of Taken’s embedding theorem [24], [25], [26], [27]. This article provides additional analytical insight into the general computation performance of delay-based reservoir computers and with it introduces new possibilities to optimize reservoirs.

Originally, reservoir computing is performed with a network of nonlinear nodes, which projects the input information into a high-dimensional phase space, allowing a linear fit to regress or to linearly separate features [1]. In time-delayed reservoir computing, a single-dynamical node with delayed feedback is employed as a reservoir instead of the network [28]. The time-multiplexing procedure allows for such a single-element system to implement a recurrent ring network [28], [29], [30], [31], see Fig. 1. The absence of the need for a large number of nonlinear elements significantly reduces the complexity of the reservoir hardware implementation. Promising realizations with a single-delay-based reservoir [32], [33], [34], [35] give a first glimpse over the potential of this idea for, e.g., time-series predictions [36], [37], [38], [39], equalization tasks on nonlinearly distorted signals [40], and fast word recognition [41]. For a general overview, we refer to [42], [43], [44], and [45].

Often reservoirs are optimized to a specific task by hyperparameter tuning, which defeats the purpose of reservoir computing as a fast trainable machine learning scheme. Dambre et al. [46] introduced a task-independent quantification of a reservoir computer, building on the memory capacity notion already introduced in [1], whereas a high memory capacity pinpoints to generally well-performing reservoirs.

In this article, we provide an analytical tool for finding promising reservoir setups by introducing a master memory function (MMF) for delay-based reservoir computing with small inputs. The MMF allows for fast computable predictions of the linear memory capacity and it indicates that the linear memory capacity of reservoirs is similar for systems with similar linearizations.

The main idea behind our method can be outlined as follows. Consider a delay-based reservoir described by a general
This article is structured as follows. First, we will briefly revise the concept of time-delay-based reservoir computing and the concept of linear memory capacity. We will then present our main analytical result while additionally presenting an example code for an efficient evaluation of the obtained expression; the derivation is given in the Appendix. Finally, comparisons of numerically simulated reservoir computer performance with the semianalytical approach are provided. We also show in Section IV-F the applications of our results to reservoirs with an unknown dynamical model, where the parameters \( a \) and \( b \) are evaluated using the system response to external small periodic stimuli.

II. TIME-DELAY-BASED RESERVOIR COMPUTING

Reservoir computing utilizes the intrinsic abilities of dynamical systems to project the input information into a high-dimensional phase space [1]. By linearly combining the responses of the dynamical reservoir, a specific task is approximated. In the classical reservoir computing scheme, often, a so-called echo state network is used by feeding the input into a spatially extended network of nonlinear nodes. Linear regression is then applied to minimize the Euclidean distance between the output and a target. This approach is particularly resourceful for time-dependent tasks because the dynamical system that forms the reservoir acts as a memory kernel.

In the time-delay-based reservoir computing scheme [28], the spatially extended network is replaced by a single-nonlinear node with a time-delayed feedback loop. A time-multiplexing procedure with a periodic mask function \( g \) is applied to translate the input data to a temporal input signal. Similarly, the time-multiplexing procedure translates the single-temporal high-dimensional reservoir response to the spatiotemporal responses of virtual nodes. The virtual nodes play the same role as the spatial nodes in echo state networks.

A sketch of the delay-based reservoir computing setup is shown in Fig. 1(a). In the following, we will give a short overview of the quantities and notations used in this article. We also refer to our previous works [27], [48], [49], [50], [51] for a detailed explanation of how the reservoir setup is operated and task-independent memory capacities are computed.

Let us briefly remind the main ingredients of the time-multiplexed reservoir computing scheme [28], [48], [49], [50]. We apply an input vector \( \mathbf{u} \in \mathbb{R}^K \) componentwise at times \( t \in [t_{k-1}, t_k) \), \( t_k = kT \), \( k = 1, \ldots, K \), \( K \) being the number of sample points. The administration time for each input \( t_{k+1} - t_k = T \) is the same and it is called the clock cycle \( T \).

To achieve a high-dimensional response to the same input, a \( T \)-periodic mask function \( g \) multiplies the input and the resulting signal enters the system (see Figs. 1 and 2). The mask \( g \) is a piecewise-constant function on \( N_V \) intervals, each of length \( \theta = T/N_V \) corresponding to \( N_V \) virtual nodes with the corresponding constant values \( w_{\text{input}} \) called the input weights. The values of the mask function \( g \) play the same role as the input weights \( w_{\text{input}} \) in spatially extended reservoirs, with the difference that time-multiplexing distributes the weights \( w_{\text{input}} \) over time. The responses of the reservoir are collected in the
state matrix $S \in \mathbb{R}^{K \times NV}$, see Fig. 3. The elements of the
state matrix are $[S]_{kn} = \hat{s}(kT + n\theta)$ with $n = 1, \ldots, NV$,
and $k = 1, \ldots, K$, where $\hat{s}(kT + n\theta) \in \mathbb{R}$ is the state of
the dynamical element of the reservoir at time $(kT + n\theta)$ shifted
by the mean over all clock cycles $\hat{s}(kT + n\theta) = \hat{s}(kT + n\theta) − \langle \hat{s}(T + n\theta) \rangle$, see [46]. The average $\langle \hat{s}(T + n\theta) \rangle$ can be understood as the averaging over the column elements,
where the $\cdot$ denotes the variable over which the averaging is
done.

A linear combination of the state matrix is given by $S w_{\text{reg}}$, where $w_{\text{reg}} \in \mathbb{R}^{NV}$ is a vector of weights. Such a combination is trained by ridge regression, i.e., the least square approximation to some target vector $\hat{y}$

$$\arg \min_{w_{\text{reg}}} \left[ \| S w_{\text{reg}} - \hat{y} \|_2^2 + \lambda \| w_{\text{reg}} \|_2^2 \right]$$

where $\| \cdot \|_2$ is the Euclidean norm, and $\lambda$ is a Tikhonov regularization parameter. The solution to this problem is

$$w_{\text{reg}} = (S^T S + \lambda I)^{-1} S^T \hat{y}.$$  

In the case of invertible $S^T S$, the matrix $(S^T S)^{-1} S^T$ is the Moore–Penrose pseudoinverse. We set $\lambda = 10^{-6}$, $\max_x(S)$, where $\max_x(S)$ is the largest state response in the state matrix $S$. To quantify the system’s performance, we use the capacity $C_{\hat{y}}$ (see [46], [49]) to approximate a specific task which is given by

$$C_{\hat{y}} = 1 - \text{NRMSE}$$  

where NRMSE [52] is the normalized root-mean-square error between the approximation $y = S w$ and the target $\hat{y}$

$$\text{NRMSE} = \left( \frac{\sum_{k=1}^{K} (\hat{y}_k - y_k)^2}{\sum_{k=1}^{K} (y_k - \bar{y})^2} \right)^{1/2}$$

where $\text{var}(\hat{y})$ is the variance of the target values $\hat{y} = (\hat{y}_1, \ldots, \hat{y}_K)$.

III. RESERVOIR COMPUTERS AND THEIR QUANTIFICATION

Here, we introduce the linear memory capacity as a quanti-
tative measure for the memory kernel of a dynamical system
and give an overview of all used reservoir computer systems
in this article.

$$\hat{y}_l = \{ \ldots, u_{-3-l}, u_{-2-l}, u_{-1-l} \}$$

which is the linear recall $l$ steps into the past. Formally, one
considers an infinitely long sequence $K \to \infty$. To approxi-
mate it numerically, we use $K = 75\,000$.  

A. Memory Capacity

The central task-independent quantification was introduced by Jaeger [1] and refined by Dambre et al. [46], which yields that the computational capability of a reservoir system can be quantified via an orthonormal set of basic functions on a sequence of inputs. Here, we give a recap of the quantities introduced in [27], [49], [50], and [51] and focus on the linear memory capacity.

In particular, the capacity to fulfill a specific task is given by

$$C_{\hat{y}} = \frac{\hat{y}^T S (S^T S + \lambda I)^{-1} S^T \hat{y}}{\| \hat{y} \|^2} = \frac{\hat{y}^T S \hat{y}}{\| \hat{y} \|^2}$$

which can be derived from (3) (see [46], [49]). The capacity equals 1 if the reservoir computer computes the task perfectly, and thus, $y = \hat{y}$; and it equals $C = 0$ if the prediction is not correlated with the target. In between 0 and 1 if it is partially capable of fulfilling the task. To quantify the system’s capability for approximating linear recalls of inputs, an input sequence $\{u\} = \{u_{-3}, \ldots, u_{-2}, u_{-1}\}$ is applied, where $u_k$ are uniformly distributed random numbers, independent and identically drawn in $[-1, 1]$. With the input sequence $\{u\}$ of random numbers, the reservoir response is collected in the state matrix $S$.

To describe a linear recall task of $l$ steps into the past, the target vector $\hat{y}$ is defined as
where $C_l = C_k$ is the capacity of the $l$th recall into the past. This quantification is task-independent, and thus, implications for specific applications cannot be given. Different tasks may need different specific capacities. The measure $MC$ thus only gives a hint for well-performing reservoirs in the context of using the full scope of the given reservoirs, rather than a direct task-specific estimate. We have to point out, that the linear-nonlinear tradeoff is a well-known effect [46], [53], and thus, a system with high linear memory capacity can yield a low nonlinear transformation capability. Nevertheless, we believe predicting a well-performing linear memory kernel reservoir is beneficial for a general reservoir computer setup, as higher nonlinear memory transformation can be utilized by adding additional reservoir systems with increased perturbations.

**B. Reservoir Systems**

The delay-based reservoir computer systems used are as follows.

1) **Stuart–Landau Oscillator** [54], [55] With Delayed Feedback:

$$\frac{ds(t)}{dt} = (p_{SL} + \eta I(t) + \gamma_{SL}(s^2(t))s(t) + \kappa s(t - \tau)).$$  (9)

Here, $s(t)$ describes the real-valued amplitude of the system, $\kappa$ is the feedback strength, $\tau$ the delay time, $p_{SL}$ is a system parameter, and $\eta$ the input strength of the information fed into the system.

For $p_{SL} + \kappa > 0$ and $\eta = 0$, (9) has only the trivial equilibrium $s^* = 0$, and for $p_{SL} + \kappa < 0$, additionally, the nontrivial equilibrium $(s^*)^2 = -(p_{SL} + \kappa)/\gamma$ exist, which appear in a pitchfork bifurcation at $p_{SL} + \kappa = 0$. The linearization at the nontrivial equilibria (taking into account the input term) reads

$$\frac{\delta s(t)}{dt} = a \delta s(t) + b \delta s(t - \tau) + c I(t)$$  (10)

where $a = -2p_{SL} - 3\kappa = p_{SL} + 3\gamma(A^*)^2$, $b = \kappa$, and $c = \eta s^*$.

2) **Mackey–Glass System** [56]:

$$\frac{ds(t)}{dt} = (p_{MG} + \eta I(t))s(t) + \frac{as(t - \tau)}{1 + s^2(t - \tau)}$$  (11)

where $s(t)$ is the dynamical variable, $s(t - \tau)$ is the delayed variable, and $p_{MG}$, $q$, and $a$ are control parameters. The reservoir input is fed into the system via the term $\eta I(t)$. We set $q = 1$, for which the system possesses a stable nontrivial equilibrium $s^* = -(p_{MG} + a)/p_{MG}$ (for $\eta = 0$). The corresponding linearization at this equilibrium is (10) with $a = p_{MG}, b = a/(1 + s^*)^2$, and $c = \eta s^*$.

**C. Overview of Used Parameters**

An overview of all used parameters is given in Table I. Table I is divided into the sections of the first parameter appearance.
IV. RESULTS

A. Analytic Description of Memory Capacity

From (6), we see that the capacity to approximate a specific input is given by the inverse of the covariance matrix \( (S^T S)^{-1} \) (corrected by \( \lambda I \)), also called the concentration matrix [59], and the matrix multiplication of the state matrix and the target \( S^T y \). Thus, it is necessary to derive the state matrix \( S \) from the responses to the small perturbations of the system. This has already been done for 1-D reservoirs with \( \tau = T \) by an Euler step scheme [17], [47], and for 1-D reservoirs with \( \tau \neq T \) for special cases of differential equations [48].

We would like to extend this knowledge by analyzing arbitrary single-variable systems and \( \tau \neq T \). We assume the node distance \( \theta \) to be small and \( \tau = \nu \theta \), with \( \nu \in \mathbb{N}^+ \). We also assume the operation point of the reservoir to be a stable equilibrium. We will exemplarily validate our approach does not require calculating the reservoir, and it can be done by two nested loops \( r \) and \( i \). We do this until all entries in a row \( q \) are below a given threshold \( \epsilon \) for \( i + r = q \) for \( i, r \in \mathbb{N} \) (see Fig. 7). The threshold \( \epsilon \) ensures that we cut unnecessary terms smaller than the regularization parameter \( \lambda \). A third loop \( n \) goes over all virtual nodes \( N_V \), adding the result \((r+i)/p^m\) multiplied with the corresponding input weights \( w_{m+i+r} \) mod \( N_V \) to all corresponding entries \( \tilde{S}_{[(r+i+r)/N_V]} \), that and thus, lie in the same input interval \( l \). See Appendix A for more information.

To compute the modified state matrix \( \tilde{S} \) is given in the following, where \( [y] \) is the floor function rounding \( y \) down to the greatest integer less than or equal to \( y \), get \( \text{BinomialTerm}(i,l,p) \) returns \( \binom{i}{l} p^m \) and \( w_{\text{input}} \) is the mask weight vector of length \( N_V \). The implemented C++ code can be found in a linked GitHub repository and can be used to calculate the memory capacity for given \( a, b, \nu, \theta, \) and \( N_V \).

Algorithm 1 Calculate Modified State Matrix

\[
\text{State} \in \mathbb{R}^{N_V \times K};
\]

\[
\text{for } \text{RowsInPascalsTriangle} < \text{maxRow do}
\]

\[
\text{for } \text{ColumnsInRow} < \text{RowNumber} + 1 \text{ do}
\]

\[
\text{for } \text{VirtualNeuron} < \text{NumberVirtualNeurons do}
\]

\[
\text{n} = \text{VirtualNeuron};
\]

\[
i = \text{RowsInPascalsTriangle};
\]

\[
r = \text{ColumnsInRow};
\]

\[
\text{State} \left( \left( \left( n + i + r \right) / N_V \right), n \right) \text{ +=}
\]

\[
\text{getBinomialTerm}(i, l, p) \cdot w_{n+i+r} \text{ mod } N_V;
\]

\[
\text{if } \text{getBinomialTerm}(i, r, p) < \sigma \text{ for all ColumnsInRow then}
\]

\[
\text{return State;}
\]

\[
\text{end for}
\]

\[
\text{end for}
\]

\[
\text{end for}
\]

\[
\text{end for}
\]

\[
\text{end for}
\]

B. Efficient Numerical Evaluation of the Memory Capacity and the Modified State Matrix

The obtained approximations of the modified state matrix (13) and memory capacity function (15) allow for efficient numerical evaluation. For this, we propose the following scheme, which we also show as a pseudocode in Algorithm 1.

First, we iterate over all entries of a modified Pascal’s triangle (see Fig. 7 in the appendix for more information on that), which can be done by two nested loops \( r \) and \( i \). We do this until all entries in a row \( q \) are below a given threshold \( \epsilon \) for \( i + r = q \) for \( i, r \in \mathbb{N} \) (see Fig. 7). The threshold \( \epsilon \) ensures that we cut unnecessary terms smaller than the regularization parameter \( \lambda \). A third loop \( n \) goes over all virtual nodes \( N_V \), adding the result \((r+i)/p^m\) multiplied with the corresponding input weights \( w_{m+i+r} \) mod \( N_V \) to all corresponding entries \( \tilde{S}_{[(r+i+r)/N_V]} \), that and thus, lie in the same input interval \( l \). See Appendix A for more information.

The algorithm to compute the modified state matrix \( \tilde{S} \) is given in the following, where \( [y] \) is the floor function rounding \( y \) down to the greatest integer less than or equal to \( y \), get \( \text{BinomialTerm}(i,l,p) \) returns \( \binom{i}{l} p^m \) and \( w_{\text{input}} \) is the mask weight vector of length \( N_V \). The implemented C++ code can be found in a linked GitHub repository and can be used to calculate the memory capacity for given \( a, b, \nu, \theta \), and \( N_V \).

C. Direct Simulation of the Reservoir and Memory Capacity

Simulations have been performed in standard C++. For linear algebra calculations, the linear algebra library “Armadillo” [60] was used. To numerically integrate the delay-differential equations, a Runge–Kutta fourth-order method was applied, with integration step \( \Delta t = 0.01 \) in dimensionless
time units. First, the system is simulated without reservoir inputs, thus letting transients decay. After that, a buffer time of 10,000 inputs was applied (this is excluded from the training process). In the training process, $K = 75,000$ inputs were used to have sufficient statistics. Afterward, the memory capacities $C_l$ of linear recalls were calculated with (6), whereby a testing phase is not necessary. The linear memory capacity $MC$ was calculated by summing the obtained capacities $C_l$. For the piecewise-constant $T$-periodic mask function $g(t)$ independent and identically distributed random numbers between $[0, 1]$ were used.

For all simulations, the input strength $\eta$ was fixed to $10^{-3}$. The small input strength was used to guarantee linear answers of the reservoir and, hence, the relevance of the approximation. In the appendix we show a detailed analysis for increased input strength.

D. Comparison of MMF and Direct Numeric Calculations of the Memory Capacity

In this section, we illustrate the MMF’s effectiveness. First, we show that the MMF provides a very good approximation of MC using the reservoir given by (9). The analytical approximation of the MC works quite well as long as $\theta$ is small enough compared with the linear answer time scale of the system, given by the largest Lyapunov exponent of the local dynamics. This is fulfilled for typical reservoir computing setups, as one would, otherwise, lose computation speed. In the second part, we show how MMF provides a universal, system-independent characteristic. For this, we compare MMF with the memory capacities of different reservoirs. Each particular reservoir realization is described by one parameter combination of linearization. In the last part, we describe how MMF can be computed for reservoirs with the unknown dynamical rule. For this, the parameters $a$ and $b$ of the linearization are measured from the system’s response to a small periodic input.

Fig. 4 shows the memory recall capacity $C_l$ obtained from direct simulations and via the MMF for four different cases of the Stuart–Landau system, given by (9). The exact parameters are given in the caption of Fig. 4. The directly simulated results are shown by blue solid lines and blue markers, whereby green dashed lines and green markers show the MMF. For a small virtual node distance $\theta = 0.5$ in Fig. 4(a) and (c), the MMF predicts the linear memory capacity very accurately. For a higher value of $\theta = 1.6$ [Fig. 4(b) and (d)], the accuracy drops, though the results are still accurate for qualitative predictions and describe the general trend of the system’s memory capacity.

The scans in Fig. 4(c) and (d) were done with a higher delay time $\tau = 3.06 T$, which induces memory gaps [50]. Even though the memory capacity has a complex dependency on $l$ at these parameter values, the prediction for the two different virtual node distances $\theta = 0.5$ and $\theta = 1.6$ is still accurate.

A 2-D parameter scan shown in Appendix E indicates that the predictions for the MC made by the MMF work for all system parameters investigated there. Thus, the predictive power of the new scheme is very promising.

Comparing the computation speed of the classical numerical integration and the new proposed scheme shows an increase in two to three orders of magnitude, depending on the operation point, the number of training steps $K$, and the value of the clock cycle $T$. A higher clock cycle $T$ and more training steps $K$ increase the simulation time for the direct numerical integration, whereas the new proposed scheme is independent of that. If the operation point is close to a bifurcation, the convergence of the new proposed scheme is slower, increasing the computation time needed. See Appendix B for a plot of the computation speed comparison in dependence of the linearization parameter $b$ close to a bifurcation. Still, even very close to the bifurcation line, the computation speed is significantly higher (with a factor of about 100) making the MMF a valuable tool.

We also performed an analysis of the valid approximation range for the input strength $\eta$ and virtual node distance $\theta$. Both results are shown in Appendix C. Fig. 9 shows the results for the Stuart–Landau delay-based reservoir computer for different input strengths $\eta$, for which $\eta$ of up to $10^{-2}$ yield comparable good predictions of the MC via the MMF. Higher values start to induce nonlinear answers in the system, thus inducing the well-known linear–nonlinear tradeoff.

Fig. 10 in Appendix C shows the results for the same system over the virtual node distance $\theta$, ranging from $\theta = 0.05$ up to $\theta = 5$. The system was simulated with different $p_{SL}$, thus changing the linear answer time-scale $a$ from $a = -0.11$ up to $a = -0.29$. The faster the system reacts the more inaccurate the predictions of the MMF are for higher values of $\theta$. An explanation for this can be given via the assumption used in Section A. Because $\theta$ is always small in reservoir computing, we assumed a constant value of $s(t)$ on one $\theta$-interval in (27). This approximation only holds for $\theta$, which is small relative to the linear answer time-scale $a$. The rule of thumb used in [28] is given by $\theta = 0.2a$. Looking at Fig. 10, we see that our approximation holds for values of up to $\theta = 3a$ for a relative error of $\Delta MC \approx 0.1$. We can, thus, conclude that
the assumption of a constant value on one $\theta$ interval for a delay-based reservoir computer is justified.

E. Universality

An exciting result that follows from the MMF concept is the possibility to generalize to any 1-D reservoir with one time-delay. Every 1-D reservoir with one time-delay that has the same linearization yields the same linear memory capacity. To illustrate this, we compare the Stuart–Landau reservoir system given by (9) and the Mackey–Glass reservoir system given by (11). The inset of Fig. 5 shows the capacity to recall the $l$th step into the past $C_l$ as a function of $l$ for the Stuart–Landau (blue), the Mackey–Glass (red), and the MMF given by (16) (green). Both systems are operated such that their linearization parameters $a$ and $b$ are equal. It can be seen that in this case the $C_l$’s have the same values.

This underlines that it is enough to compute the linearization parameters $a$ and $b$ to predict the MC of a single-variable delay-based reservoir computer. The color plot in Fig. 5 shows the MMF given by (15) for different parameter values $a$ and $b$. A well-performing operation point seems to be the edge to instability, agreeing with the known rule of thumb from the literature [61], [62].

It, thus, follows that analyzing the Jacobian [linearization given by (10)] for fixed delay $\tau$, virtual node distance $\theta$, and the number of virtual node $N_v$ is sufficient to predict the linear memory capacity of a single-variable time-delay-based reservoir computer, and this memory capacity is given by the MMF via (15) and (16).

F. Systems With Unknown Dynamics—Small Signal Response Approach

In this chapter, we show an experimentally accessible approach for measuring the parameters $a$ and $b$ for a delay system whose dynamical equations of motion are not known and which can be described by a single variable. The linearized dynamical system was introduced in (10). To measure $a$ and $b$, we perturb the system with a harmonic periodic signal $I(t) = I_0 e^{i\omega t}$. When this signal is small, we can consider the perturbed linearized system

$$\frac{d\delta s(t)}{dt} = a\delta s(t) + b\delta s(t-\tau) + cI_0 e^{i\omega t}$$  \hspace{2cm} (17)$$

where the complex form is chosen out of convenience. Due to linearity, the real solution is obtained simply by taking the real part. We consider the case of real $a$ and $b$, which always holds when the reservoir variables are real.

Since the homogeneous solution decays to the stable equilibrium (we assume its exponential stability), the solution of (17) converges to the particular solution, given by

$$\delta s_a(t) = cI_0 H^{-1}(\omega)e^{i\omega t}$$  \hspace{2cm} (18)$$

with $H^{-1}(\omega) = i\omega - a - be^{-i\omega \tau}$. The ratio of the output to the input amplitude equals the transfer function

$$\frac{|\text{Output}|}{|\text{Input}|} = \frac{|\delta s_a(t)|}{|I_0|} = |H^{-1}(\omega)|$$  \hspace{2cm} (19)$$

where $|H(\omega)|$ can be measured.

To determine the parameters $a$ and $b$, it is sufficient to measure the transfer function at two frequencies, for example, at $\omega_R = 2\pi/\tau$ and $\omega_A = \pi/\tau$. The first frequency is resonant to the delay while the second is in “antiphase” to the delay $\tau$. We assume $\tau$ to be known, even though an extension of the scheme to unknown $\tau$ is possible. It holds

$$F(\omega_R) := |H(\omega_R)|^2 = \omega_R^2 + (a + b)^2$$
$$F(\omega_A) := |H(\omega_A)|^2 = \omega_A^2 + (a - b)^2.$$  \hspace{2cm} (20)$$

From the above, we can obtain the values for $a$ and $b$

$$a = -\frac{1}{2} \left( \sqrt{F(\omega_R) - \omega_R^2} + \sqrt{F(\omega_A) - \omega_A^2} \right)$$
$$b = -\frac{1}{2} \left( \sqrt{F(\omega_A) - \omega_A^2} - \sqrt{F(\omega_R) - \omega_R^2} \right)$$  \hspace{2cm} (21)$$

where the values of $F(\omega_A)$ and $F(\omega_R)$ can be obtained experimentally or numerically by perturbing and measuring the response of the reservoir.

We remark that the choices of the resonant and antiphase perturbation frequencies are convenient but not unique. Clearly, one can perturb at other frequencies to obtain (20) and (21).

The measured values of the parameters $a$ and $b$ for a reservoir with unknown dynamics can be then simply used in MMF to determine the linear memory capacities.

To demonstrate the experimental procedure of measuring $a$ and $b$ via a small signal response, we use the Ikeda-system introduced in Section III-B in (12) with $x_0 = 1.5$, $\tau = 72$, and varying nonlinearity parameter $\mu$ ranging from $\mu_{\text{min}} = 0.1$ up to $\mu_{\text{max}} = 1.2$. In this range, the system has one stable fixed-point solution. We choose $I(t) = \cos(\omega t)$. With $\omega_R = 2\pi/\tau$, $\omega_A = \pi/\tau$, and $\eta = 0.01$. We simulate the systems response and measure $|s_{a,R}(t)| = \max |s(t) - s^*|$, thus $|s_{a,R/A}(t)|$ is the induced linear response oscillation amplitude superimposed on the fixed point $s^*$ for the in-phase $|s_{a,R}(t)|$ or antiphase $|s_{a,A}(t)|$ case.
The results are shown in Fig. 6(a). The red curve shows the measured values of \(a\) and \(b\) over the Ikeda parameter \(\mu\). We additionally calculated the linearization parameters via the linearization that was derived in Section III-B. The result is shown as a green solid line in Fig. 6(a). We see that the measured and the true values are nearly identical, and thus, the measuring method is working as expected. We then used one example parameter pair shown as a black cross in Fig. 6(a) and simulated the full Ikeda-system as a delay-based reservoir computer at this parameter setup, numerically evaluated the linear memory capacities \(C_l\), and plotted them over the recall steps \(l\) in Fig. 6(b) as a blue curve. Subsequently, the MMF was used to calculate \(C_l\) for the measured \(a = -1\) and \(b \approx -0.6169\) over the same recall steps \(l\). The results are shown in Fig. 6(b) as a green dashed line. The fully numerically simulated and calculated memory capacities \(C_l\) and the MMF values obtained after measuring \(a\) and \(b\) show very well the potential of the MMF.

V. CONCLUSION

We have developed a simple and fast method for calculating the linear memory capacity for time-delay-based reservoirs with single-variable dynamics, which we call the MMF. Our results can be used to predict reservoir computing setups with high linear memory capacities independent of the specific reservoir computer. We have shown universality properties for our MMF for any single-variable delay-based reservoir computer with small inputs, i.e., our approach predicts the linear memory capacity of the set of all single-variable delay-based reservoirs with small inputs. We additionally provided a method for measuring the linearization parameters \(a\) and \(b\) for dynamical systems with unknown models, therefore, enabling the possibility for experimental evaluation of the MMF. An example approach for the Ikeda-system was given.

One of the advantages of the delay-based reservoir, which allows the introduction of the MMF, is that it contains a small number of system parameters, while the dynamics remain infinite-dimensional. In the case of a small input signal and single-variable dynamics, these are only the linearization parameters \(a\) and \(b\), the time-delay \(\tau\), the virtual node distance \(\theta\), and the number of virtual nodes \(N_v\). Thus, if the linear memory capacity is computed for all possible values of these parameters, it covers the case of all possible reservoirs. This procedure could be difficult, if not impossible, for network-based reservoirs, where the system’s parameters may include, e.g., multiple coupling weights.

APPENDIX A
DERIVATION OF THE MODIFIED STATE MATRIX AND REDUCED FORMULA FOR MEMORY CAPACITY

Consider a single-variable delay-differential equation, which describes the dynamics of the reservoir

\[
\dot{s}(t) = F(s(t), s(t-\tau), I(t))
\]

where \(I(t)\) stands for an external input. We assume that \(s^*\) is an equilibrium of this system, i.e., \(F(s^*, s^*, 0) = 0\), and \(I(t)\) is small. In the case when the dynamics of (22) takes place near the equilibrium point \(s^*\), we can introduce the perturbative ansatz \(s(t) = s^* + \delta s(t)\). Then, the linearized system for the perturbation \(\delta s(t)\) reads

\[
\dot{\delta s}(t) = a \delta s(t) + b \delta s(t-\tau) + c I(t)
\]

where \(a = \partial_1 F(s^*, s^*, 0), b = \partial_2 F(s^*, s^*, 0), \) and \(c = \partial_3 F(s^*, s^*, 0)\).

Consider \(\theta\) to be the temporal node spacing of the reservoir, which is typically \(\theta < \tau\). Then, (23) on any interval \([j\theta, (j+1)\theta]\) can be considered as the simple scalar ordinary differential equation (ODE) \(\dot{\delta s}(t) = a \delta s(t) + b \delta s(t-\tau) + c I(t)\). Moreover, according to the reservoir setup, the input is constant on this interval and equals \(I(t) = I_j\). By variation of constants formula, we obtain the solution of (23) for \(t \in [(j-1)\theta, j\theta]\)

\[
\delta s(t) = -\frac{c I_j}{a} \left( 1 - e^{a(t-t_j)} \right) + \delta s(t_{j-1}) e^{a(t-t_{j-1})} + b \int_{t_{j-1}}^{t} e^{a(t-\zeta)} \delta s(\zeta - \tau) d\zeta
\]

(24)

where \(t_{j-1} = (j-1)\theta\) is the left endpoint of the interval.

Denoting \(\delta s_j(t) = \delta s(t_{j-1} + t)\) to be the function on the interval \([t_{j-1}, t_{j-1} + \theta]\), with \(t \in (0, \theta)\), we rewrite (24) as

\[
\delta s_j(t) = -\frac{c I_j}{a} \left( 1 - e^{at} \right) + \delta s_{j-1}(\theta) e^{at} + b \int_{0}^{t} e^{a(t-\zeta)} \delta s_{j-1}(\zeta) d\zeta, \quad t \in [0, \theta]
\]

(25)

where we additionally used the relation \(\delta s_{j-1}(\theta) = \delta s_j(0)\) and \(\delta s_{j-\nu} = \delta s_j(t-\tau), \nu = \tau/\theta\). By evaluating (25) at \(t = \theta\), we obtain

\[
\delta s_j(\theta) = -\frac{c I_j}{a} (1 - p) + \delta s_{j-1}(\theta) p + b p \int_{0}^{\theta} e^{-a\zeta} \delta s_{j-1}(\zeta) d\zeta, \quad p = e^{a\theta}.
\]

(26)

Denote \(s_j := s_j(\theta) = s^* + \delta s_j(\theta)\), which is the approximation for state of the reservoir (22) at the virtual node \(s(j\theta)\).
From (26), we obtain

$$s_j = (1 + \frac{b}{a}p)(1 - p)s^* - \frac{c}{a}I_j(1 - p) + ps_{j-1} + bp \int_0^\theta e^{-a\zeta} s_{j-v}(\zeta) d\zeta. \quad (27)$$

Furthermore, we approximate the integral from (27) by assuming $s_{j-v}(\zeta) \approx s_{j-v}(\theta) = s_{j-v}$. The approximation holds, in particular, when $\theta$ is small. The obtained expression

$$s_j = \hat{s}^* + \gamma I_j + ps_{j-1} + ms_{j-v} \quad (28)$$

represents a discrete map (coupled map lattice) for approximating the state matrix $S$. Here, $\hat{s}^* = (1 - (b/a)p)(1 - p)s^*$, $m = -(b/a)(1 - p)$, and $\gamma = -(c/a)(1 - p)$. If considering it as a corresponding network with the nodes $s_j$, see [29], [31], [63], we see that the node $s_j$ is coupled with the two nodes $s_{j-1}$ and $s_{j-v}$ in a feed-forward manner with the coupling weights $p$ and $m$, respectively. The schematic representation of such a coupling structure leads to Pascal’s triangle shown in Fig. 7. The first row of Pascal’s triangle from Fig. 7 shows the dependence on $I_j$, which is simply the multiplication by $\gamma$. In the second row, the contributions of $I_{j-1}$ and $I_{j-v}$ are shown. To obtain these dependencies explicitly, we insert $s_{j-1}$ and $s_{j-v}$ recursively in (28)

$$s_j = (1 + p + m)\hat{s}^* + \gamma (I_j + pI_{j-1} + mI_{j-v}) + p^2s_{j-2} + 2ps_{j-v-1} + m^2s_{j-2v} \quad (29)$$

that is, we obtain the terms $\gamma I_{j-1}$ and $\gamma mI_{j-v}$. To build up further intuition about the dependence of the state matrix on the input, we show here the third level by substituting recursively $s_{j-2}$, $s_{j-v-1}$, and $s_{j-2v}$ into (29)

$$s_j = (1 + p + p^2 + 2pm + m^2)\hat{s}^* + \gamma (I_j + pI_{j-1} + mI_{j-v} + p^2I_{j-2} + 2pmI_{j-v-1} + m^2I_{j-2v}) + p^2s_{j-3} + 3p^2ms_{j-v-2} + 3pm^2s_{j-v-2} + m^3s_{j-3v}. \quad (30)$$

To obtain a general recursive formula, we need to split the index into the appearing terms as $j - i - kv$, where $k$ corresponds to the delayed (“right,” $r$) and $i$ to the “left” ($p$) connections in the coupling network in Fig. 7

$$s_j = A_1s^* + \gamma \sum_{i, r=0}^\infty (r + i) p^rm^r I_{j-i-rv} \quad (31)$$

where $A_1$ is a constant depending only on $p$ and $m$. For an infinitely long input sequence, the sum in (31) goes for all $i$, $r$ from 0 to $\infty$. Practically, the sum is considered for the available data $I_j$. As a result, the reservoir states $s_j$ are composed of a linear combination of the inputs with corresponding coefficients given in (31). We can drop the $j$-independent constant term $A_1s^*$, because it only induces a constant shift of the state matrix $S$, and thus, can be absorbed in either a bias term or a centering preprocessing, one of the two is always done when calculating the memory capacity.

The elements of the state matrix $S$ used in the reservoir computing setup are

$$[S]_{kn} = \hat{s}_{KNv+n} - s_{KNv+n} \quad (32)$$

where $s_{KNv+n}$ is the average over the input intervals $\langle s_{Nv+n} \rangle = (1/K) \sum_{k=1}^K s_{KNv+n}$. Here and later, the dot denotes the index, over which the averaging is performed. Taking into account (31), we obtain

$$\hat{s}_{KNv+n} = \gamma \sum_{i, r=0}^\infty (r + i) p^rm^r \times (I_{KNv+n-i-rv} - \langle I_{Nv+n-i-rv} \rangle) \quad (33)$$

The input $I_k$ of the reservoir computer is given by the discrete input sequence $u$ multiplied by the input weights

$$I_k = u[k/Nv]I_{k \mod Nv}. \quad (34)$$

Therefore, we obtain

$$\langle I_{Nv+n-i-rv} \rangle = u[n-i-rv \mod Nv] \langle I_{n+(n-i-rv)/Nv} \rangle \approx 0 \quad (35)$$

since $u$ has zero mean. Hence, we have the elements of the state matrix

$$\hat{s}_{KNv+n} = \gamma \sum_{i, r=0}^\infty (r + i) p^rm^r I_{KNv+n-i-rv}. \quad (36)$$

Correspondingly, the elements of the covariance matrix $S^T S$ from (6) are

$$[S^T S]_{kn} = \sum_k [S]_{kn} [S]_{kn} = \sum_k \hat{s}_{KNv+n} \hat{s}_{Knv+n'} \quad (37)$$

and they describe the covariance of the virtual node $n$ with virtual node $n'$ over all clock cycles $k$.

By substituting (34) into (35), we obtain

$$[S^T S]_{kn} = \gamma^2 \sum_k \left( \sum_{r, i} (r + i) p^rm^r I_{Knv+n-i-rv} \right) \times \left( \sum_{r', i'} (r' + i') p^{r'}m'^{r'} I_{Knv+n'-i'-rv'} \right). \quad (38)$$
One can show that the elements from $[S^T S]_{nn'}$ containing mixed terms of the form $u_i u_j$, $i \neq j$ can be approximated by zero since the random variable $u_j$ is independently distributed with zero mean. The only nonzero second moment, which matters in $[S^T S]_{nn'}$, is the mean square

$$\sum_{k=0}^{\infty} u_k^2 = \frac{1}{3}.$$  

(37)

Hence, for further simplification of $[S^T S]_{nn'}$, we keep only terms of the form $u_i^2$. The following calculations formalize these arguments, equation (38), as shown at the bottom of the page, where the second summation range ($*$) is taken over all values of $r, i, r', i'$ such that $n + l N_V \leq i + r v < n + (l + 1) N_V$ and $n' + l N_V \leq i' + r' v < n' + (l + 1) N_V$. The obtained expression (38) does not depend on the sequence $u_k$, and hence, provides a significant simplification for calculating the covariance matrix. We may further notice that the same covariance (38) can be obtained by defining the modified state matrix $\tilde{S} = \tilde{s}_n$, where $i$ is the $l$th interval of the shifted input (the $l$th recall) and $n$ the $n$th virtual node. $\tilde{s}_n$ is given by the sum over all combinations $i, r$ that fall into the same shifted input interval $j$, that is,

$$\tilde{s}_n = \frac{\gamma}{\sqrt{3}} \sum_{n+\mathcal{L} V \leq i + r v} \binom{r + i}{i} p^i m^r u_{(n-i-r v) \bmod N_V}.$$  

(39)

This is our main result because (39) defines the modified state matrix $\tilde{S}$ from which all capacities $C_l$ are derivable. More specifically, we have shown

$$S^T S \approx S^T \tilde{S}.$$  

(40)

Furthermore, for the $l$th recall, where the target is the shifted input sequence $\tilde{y} = \{u_{k-l}\}_{k=1}^{\infty}$, we have, equation (41), as shown at the bottom of the page, therefore, it holds

$$S^T \tilde{y}_i \approx \frac{1}{\sqrt{3}} \tilde{s}_n$$

where $\tilde{S}_l$ is the $l$th row of $\tilde{S}$. Furthermore, we notice that

$$\hat{y}^T_l S = (S^T \hat{y}_i) \approx \frac{1}{\sqrt{3}} \tilde{S}_l^T$$

and $\|\hat{y}_i\|^2 \approx 1/3$. As a result, taking into account the definition of the memory capacity (6), we obtain the approximation for the capacity of the $l$th recall $C_l$ by

$$C_l \approx \tilde{S}_l^T (S^T S + \lambda I)^{-1} \tilde{s}_n.$$  

(42)

The results can be understood in such a way that we constructed the modified state matrix $\tilde{S}$, such that every column has entries in the statistical direction of the $l$th shifted input recall.

The full linear memory capacity is then given by the trace

$$MC_{MMF} = \text{tr} \left( S^T \tilde{S} \tilde{S}^T \right).$$  

(43)

\section*{Appendix B: Computation Time}

To compare the computation speed of the full numerically simulated differential equation and our new analytic approach, we simulated both systems. The full system with a timestep of $dt = 0.01$, buffer samples of 10000, i.e., that 10000 clock cycles were simulated and discarded, and 50 000 training
samples to get high accuracy on the memory capacity. The analytic program was calculated until all values in a row in Pascal’s triangle were below $10^{-6} \cdot \max(S)$. We compared the simulation speeds of both approaches on the same hardware on a parameter line scan of the linearization parameter $b$, scanning from values close to the bifurcation value $b_{\text{bif}}$ in which the linearized system destabilizes up to values of about 0.1 greater than the bifurcation value $b_{\text{bif}}$. We show the percentage of the simulation time for the analytic approach $t_{\text{anal}}$ in comparison with the simulation time $t_{\text{sim}}$, i.e., $t_{\text{anal}} \div t_{\text{sim}}$ in Fig. 8. We see that close to the bifurcation the analytic approach increases in computation time. This comes from the fact that the convergence of Pascal’s triangle close to the bifurcation is slower. Still, the simulation time is at maximum 4% of the fully simulated system, showing at least a 25-fold increase in computation speed.

APPENDIX C
RANGE OF APPROXIMATION

We would like to show the range of approximation for the new analytic approach for both the input strength $\eta$ and the virtual node distance $\theta$. For the input strength $\eta$, we compute the memory capacity $MC_{\text{sim}}$ of the fully simulated system $MC_{\text{sim}}$ and the analytic approach $MC_{\text{anal}}$ and show the relative memory capacity of the analytic approach to the full system, i.e., $MC_{\text{anal}} \div MC_{\text{sim}}$. The results are shown in Fig. 9 plotted over the input strength $\eta$ for six orders of magnitude. A result close to 1 indicates a good agreement between the simulation and the analytic approach. For high input strengths, starting at around $\eta = 10^{-2}$, the analytic approach overestimates the real memory capacity, because high values of $\eta$ induce nonlinear answers in the system, and thus, increase the nonlinear transformations of the reservoir in exchange for linear memory, see [46], [49], [50] for more information on that effect. This can be seen when the MC of the fully simulated system $MC_{\text{sim}}$ is shown, plotted for the second y-axis in red.

As a second quantity, we plot the relative difference of the MC, that is,

$$\Delta MC = \frac{|MC_{\text{MMF}} - MC_{\text{direct}}|}{MC_{\text{direct}}}$$

over the virtual node distance $\theta$ for different values of the Stuart–Landau control parameter $p_{\text{SL}}$, where $MC_{\text{MMF}}$ is the prediction of the MC for the MMF and $MC_{\text{direct}}$ is the true MC computed via a fully numerical simulation. The results are shown in Fig. 10. The different colored graphs depict different operation points of the Stuart–Landau delay-based reservoir computer $p_{\text{SL}}$ ranging from $p_{\text{SL}} = -0.005$ up to $p_{\text{SL}} = -0.095$, which corresponds to linearization parameters of $a = -0.11$ up to $a = -0.29$, respectively. With increasing $\theta$ the prediction of the MMF becomes less accurate. This effect is intensified if the linearized local dynamics are faster, i.e., for more negative $a$. Faster local dynamics result in faster changing dynamics, and thus, the assumption of a constant...
\[ \delta s(t) \text{ on one virtual node interval } \theta \text{ becomes increasingly inaccurate.} \]

**APPENDIX D**

\[ \chi_k^2 \text{ Estimation} \]

We give a short insight into the \( \chi_k^2 \) estimation introduced in [46]. When calculating capacities \( C_1 \), all below a fixed value \( r^* \) were excluded because of finite statistics, where \( r^* \) is given by the following relation. CDF(\( \chi^2(N_V, r^*) \)) is the cumulative distribution function of the \( \chi^2 \) function, and \( r^* \) is chosen such that \( 1 - \text{CDF}(\chi^2(N_V, r^*)) \) yields a probability \( p_{r^2} = 10^{-6} \), i.e., the probability of a capacity having a value greater than \( r^* \) even though with infinite statistics (\( K \to \infty \)), it would have a value less than \( r^* \). \( \chi^2 \) is the probability density function of the sum of squared independent, standard normal random variables

\[ \chi_k^2 = \sum_{i=1}^{k} Z_i^2. \]  

See [46] for more information.

**APPENDIX E**

**Broader Parameter Range Check**

A two-parameter characterization of the memory capacity of the Stuart–Landau system (9) is shown in Fig. 11. The parameter space is spanned by the pump \( p_{\text{PSL}} \) and the feedback rate \( \kappa \). Fig. 11(a) shows the linear memory capacity, while Fig. 11(b) shows the relative difference \( \Delta MC \) of the MMF and the direct numerics [see (44)]. Small relative differences of up to 0.08 are seen for the simulations presented here for \( \theta = 1 \). One has to remember that reservoir computing is usually done with very small \( \theta \). The work [1] introduced a rough estimate of the optimal value for \( \theta \) as \( \theta \approx 0.2 \lambda_{\text{ans}} \), where \( \lambda_{\text{ans}} \) is the linear answer timescale of the system. In the case of the Stuart–Landau system, this is given by \( \lambda_{\text{ans}} = -2 p_{\text{PSL}} - 3 \kappa \). For the parameter space in Fig. 11, the highest value of the linear answer timescale \( \lambda_{\text{ans}} = -0.95 \), and thus, \( \theta \) is by a factor of five bigger than the proposed value given in [1] for optimal virtual node distance. In our approximation, we assume a constant state value on one \( \theta \)-interval, and thus, \( \theta = 1 \) is a very high value, which is one of the reasons for the deviations.

To underline that the MMF (15) gives a reasonable estimation of the memory capacity, we also calculated the 2-D correlation coefficient \( RV(Y, X) \) between the directly simulated total linear memory capacity \( MC_{\text{direct}} \) and the linear memory capacity \( MC_{\text{MMF}} \) given by MMF in the 2-D plane of the pump \( p_{\text{PSL}} \) and the feedback rate \( \kappa \) parameters. \( RV(Y, X) \) is the generalization of the squared Pearson coefficient for two dimensions and is calculated via

\[ RV(Y, X) = \frac{\text{COV}(X, Y)}{\sqrt{\text{VAR}(X) \text{VAR}(Y)}} \]  

with

\[ \Sigma_{XY} = E(X^T Y) \]  

\[ \text{COV}(X, Y) = \text{tr}(\Sigma_{XY} \Sigma_{YX}) \]  

\[ \text{VAR}(X) = \text{tr}(\Sigma_{XX}). \]

Here, \( E() \) is the expectation value, \( \Sigma_{XY} \) denotes the centered covariance matrix of the matrices \( X \) and \( Y \), \( \text{COV}(X, Y) \) denotes the trace of the matrix multiplication of \( \Sigma_{XY} \Sigma_{YX} \) and \( \text{VAR}(X) \) the trace of the matrix multiplication \( \Sigma_{XX} \). Calculating \( RV \) over the parameter range shown in Fig. 11 yields a value of \( RV(MC_{\text{direct}}, MC_{\text{MMF}}) \approx 0.99925 \). The correlation is close to the maximum of 1, allowing us to make accurate predictions of high-performing reservoirs with the MMF.

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