A lattice of combinatorial Hopf algebras,
Application to binary trees with multiplicities.

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Abstract. In a first part, we formalize the construction of combinatorial Hopf algebras from plactic-like monoids using polynomial realizations. Thanks to this construction we reveal a lattice structure on those combinatorial Hopf algebras. As an application, we construct a new combinatorial Hopf algebra on binary trees with multiplicities and use it to prove a hook length formula for those trees.

Résumé. Dans une première partie, nous formalisons la construction d’algèbres de Hopf combinatoires à partir d’une réalisation polynomiale et de monoïdes de type monoïde plaxique. Grâce à cette construction, nous mettons à jour une structure de treillis sur ces algèbres de Hopf combinatoires. Comme application, nous construisons une nouvelle algèbre de Hopf sur des arbres binaires à multiplicités et on l’utilise pour démontrer une formule des équerres sur ces arbres.

Keywords: Combinatorial Hopf algebras, monoids, polynomial realization, hook length formula, generating series, binary trees

1 Introduction

In the past decade a large amount of work in algebraic combinatorics has been done around combinatorial Hopf algebras. Many have been constructed on various combinatorial objects such as partitions (symmetric functions [Mac95]), compositions (NCSF [KL+94, MR95]), permutations (FQSym [MR95, DHT02]), set-partitions (WQSym [Hiv99, NT06]), binary trees (PBT or the Loday-Ronco Hopf algebra of planar binary trees [LR98, HNT04, HNT05]), or parking functions (PQSym [NT04, NT07]).

A powerful method to construct those algebras, called polynomial realization, is to construct the Hopf algebra as a sub algebra of a free algebra of polynomials (commutative or not) admitting certain symmetries. Beside the contruction of Combinatorial Hopf algebra, several recent papers investigate toward the formalization of combinatorial applications such as hook formulas, or seek some structure in this large zoo.

This extended abstract, reports on a work in progress which proposes to formalize the construction of Hopf algebras by polynomial realizations: starting with one of the three Hopf algebras FQSym, WQSym or PQSym realized in a free algebra, we impose some relations on the variables. Under some simple hypotheses, the result is again a Hopf algebra (Theorem [ ]). Two important examples are already known,
namely the Poirier-Reutenauer algebra of tableaux \((\text{FSym})\) obtained from the plactic monoid \([\text{LS}81]\) and the planar binary tree algebra of Loday-Ronco obtained from the sylvester monoid \([\text{LR}98, \text{HNT}04, \text{HNT}05]\).

We further observe that the construction transports the lattice structure on monoids to a lattice structure on those Hopf algebras (Theorem \(\text{3}\)). This structure was used implicitly by Giraudo for constructing the Baxter Hopf algebra from the Baxter monoid as the infimum of the sylvester monoid and its image under Schützenberger involution. The supremum of those two monoids is known as the hypoplactic monoid which gives the algebra of quasi symmetric functions \([\text{Nov}00]\).

As an application (Section \(\text{5}\)) we take the supremum of the sylvester monoid and the stalactic monoid of \([\text{HNT}08a]\). The result is a monoid on binary search trees with multiplicities leading to a Hopf algebra on binary trees with multiplicities. Interestingly, there is a hook length formula for those trees (Theorem \(\text{3}\)) and we prove it using the Hopf algebra as generating series.

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### 2 Background

In this section, we introduce some notations and three specific maps from words to words: standardization, packing, and parkization. These will be the main tool for polynomial realizations of Hopf algebras.

#### 2.1 Lattice structure on Congruences

The free monoid \(A^*\) on an alphabet \(A\) is the set of words with concatenation as multiplication. We denote by \(1\) the empty word. Recall that a monoid congruence is an equivalence relation \(\equiv\) which is left and right compatible with the product; in other words, for any monoid elements \(a, b, c, d\), if \(a \equiv b\) and \(c \equiv d\) then \(ac \equiv bd\). Starting with two congruences on can build two new congruences:

- the union \(\sim \lor \approx\) of \(\sim\) and \(\approx\) is the transitive closure of the union \(\sim\) and \(\approx\); that is \(u \equiv v\) if there exists \(u = u_0, \ldots, u_k = v\) such that for any \(i\), \(u_i \sim u_{i+1}\) or \(u_i \approx u_{i+1}\). It is the smallest congruence containing both \(\sim\) and \(\approx\);
- the intersection \(\sim \land \approx\) of \(\sim\) and \(\approx\) is defined as the relation \(\equiv\) with \(u \equiv v\) if \(u \sim v\) and \(u \approx v\).

#### 2.2 Some \(\varphi\)-maps

Throughout this paper we construct Hopf algebras from the equivalence classes of words given by the fibers of some map \(\varphi\) from the free monoid to itself. Our main examples are standardization and packing functions which can be defined for any totally ordered alphabet \(A\). We could easily extend these following properties to parkization \([\text{NT}04, \text{NT}07]\) if the alphabet \(A\) is well-ordered (any element has a successor).

In the following, we suppose that \(A\) is an totally ordered infinite alphabet. Most of the time we use \(A = \mathbb{N}^{>0}\) for simplicity. For \(w\) in \(A^*\), we denote by \(\text{part}(w)\) the ordered set partition of positions of \(w\).
In this section we describe how, from a \(\varphi\)-map, one can construct a Hopf algebra such as \(\text{FQSym}, \text{WQSym}\), or \(\text{PQSym}\), using two tricks: polynomial realization and alphabet doubling. Polynomial realizations are a powerful trick to construct algebras as sub-algebras of a free algebra by manipulating some polynomials having certain symmetries. Furthermore the alphabet doubling trick defines a graded algebra morphism on a free algebra which endows it with a compatible coproduct, that is a Hopf algebra structure.
3.1 $\varphi$-polynomial realization

The notion of polynomial realizations has been introduced and implicitly used in many articles of the “phalanstère de Marne-la-Vallée” (France). See e.g. [DHT02, NT06a, HNT08a]. In the following, we call alphabet $\mathfrak{A}$ an infinite and totally ordered (when appropriate, we assume furthermore that the total order admits a successor function) set of symbols all of which are of weight 1. By an abuse of language, we call the free algebra the graded algebra infinite but finite degree sum of words.

**Definition 1 (Polynomial realization):** Let $\mathfrak{A} := \bigoplus_{n \geq 0} \mathfrak{A}_n$ be a graded algebra. A polynomial realization $r$ of $\mathfrak{A}$ is a map which associates to each alphabet $\mathfrak{A}$ an injective graded algebra morphism $r_\mathfrak{A}$ from $\mathfrak{A}$ to the free non-commutative algebra $K\langle \mathfrak{A} \rangle$ such that, if $\mathfrak{A} \subset \mathfrak{B}$, then for all $x \in \mathfrak{A}$ one has $r_\mathfrak{A}(x) = r_\mathfrak{B}(x)/\mathfrak{A}$, where $r_\mathfrak{B}(x)/\mathfrak{A}$ is the sub linear combination obtained from $r_\mathfrak{B}(x)$ by keeping only those words in $\mathfrak{A}^*$. When the realization is clear from the context we write $A(\mathfrak{A}) := r_\mathfrak{A}(A)$ for short.

For a given $\varphi$, we consider the subspace $A_\varphi$ admitting the basis $(m_u)_{u \in can_\varphi}$ defined on $A_\varphi(\mathfrak{A})$:

$$r_\mathfrak{A},\varphi(m_u) = \sum_{w \in \mathfrak{A}^*; \varphi(w) = u} w.$$ (1)

The result does not depend on the alphabet. For $\varphi = std$, pack or park the linear span of $(m_u)_{u \in can,\varphi}$ is a sub-algebra of $K\langle \mathfrak{A} \rangle$.

**Example 1 (Realization of FQSym):** If $\varphi = std$ then $can_\varphi$ is in fact the set of permutations and $\mathfrak{A}_\varphi$ is the permutations Hopf algebra FQSym [DHT02, MR95]. It is realized by the $std$-polynomial realization in $K\langle \mathfrak{A} \rangle$: let $G_\sigma(\mathfrak{A}) := r_\mathfrak{A},std(G_\sigma)$ such that, for example

$$G_{132}(N^*) = 121 + 131 + 132 + 141 + 142 + 143 + \cdots + 242 + 243 + \cdots.$$

The realization is an algebra morphism: $G_\sigma(\mathfrak{A}) \cdot G_\mu(\mathfrak{A}) = r_\mathfrak{A},std(G_\sigma \times G_\mu)$ where “$\cdot$” is the classical concatenation product on words in the free algebra. For example,

$$G_{213} \times G_1 = G_{2134} + G_{2143} + G_{3142} + G_{3241}$$

which is equivalent to

$$r_{std,N^*}(G_{213} \times G_1) = G_{213}(N^*) \cdot G_1(N^*) = (212 + 213 + 214 + 215 + \cdots) \cdot (1 + 2 + 3 + 4 + 5 + \cdots) = 2121 + 2122 + 2123 + \cdots 2131 + 2132 + 2133 + \cdots + 3241 + \cdots$$

**Proposition 1:** If $span((m_u)_{u \in can_\varphi})$ is stable under the product $\times$ then it is given by:

$$m_u \times m_v = \sum_{w := u'v' \in can_\varphi; \varphi(u') = u; \varphi(v') = v} m_{w}.$$ (2)

**Remark 1:** Let $\mathfrak{A}, \mathfrak{B}$ be two totally ordered alphabets such that any element in $\mathfrak{A}$ is strictly smaller than any element of $\mathfrak{B}$. By definition we have the following isomorphisms, where $\sqcup$ denotes the disjoint union:

$$\mathfrak{A} \simeq A(\mathfrak{A}) \simeq A(\mathfrak{B}) \simeq A(\mathfrak{A} \sqcup \mathfrak{B}).$$ (3)
3.2 Alphabet doubling trick

The alphabet doubling trick [DHT02, Hiv07] is a way to define coproducts. We consider the algebra \( K\langle \mathcal{A} \sqcup \mathcal{B} \rangle \) generated by two (infinite and totally ordered) alphabets \( \mathcal{A} \) and \( \mathcal{B} \) such that the letters of \( \mathcal{A} \) are strictly smaller than the letters of \( \mathcal{B} \). The relation \( \equiv \) makes the letters of \( \mathcal{A} \) commute with those of \( \mathcal{B} \). One identifies \( K\langle \mathcal{A} \sqcup \mathcal{B} \rangle / \equiv \) with the algebra \( K\langle \mathcal{A} \rangle \otimes K\langle \mathcal{B} \rangle \). We follow here the abuse of language allowing infinite but finite degree sum. We denote by \( r_{\mathcal{A}\sqcup\mathcal{B}}(x)/\equiv \) the image of \( r_{\mathcal{A}\sqcup\mathcal{B}}(x) \) given by the canonical map from \( K\langle \mathcal{A} \sqcup \mathcal{B} \rangle \) to \( K\langle \mathcal{A} \sqcup \mathcal{B} \rangle / \equiv \). The map \( x \mapsto r_{\mathcal{A}\sqcup\mathcal{B}}(x)/\equiv \) is always an algebra morphism from \( \mathcal{A} \) to \( K\langle \mathcal{A} \rangle \otimes K\langle \mathcal{B} \rangle \). Whenever its image is included in \( \mathcal{A}(\mathcal{A}) \otimes \mathcal{A}(\mathcal{B}) \) this defines a coproduct on \( \mathcal{A} \).

**Definition 2 (Hopf polynomial realization):** A Hopf polynomial realization \( r \) of \( \mathcal{H} \) is a polynomial realization such that for all \( x \):

\[
r_{\mathcal{A}\sqcup\mathcal{B}}(x)/\equiv = (r_\mathcal{A} \otimes r_\mathcal{B})(\Delta(x)).
\]  

**Example 2 (Coproduct in FQSym):** We denote by \( G_\sigma(\mathcal{A} \sqcup \mathcal{B}) \) the \( \text{std} \)-polynomial realization of the FQSym element indexed by \( \sigma \) in the algebra \( K\langle \mathcal{A} \sqcup \mathcal{B} \rangle / \equiv \). Also we denote by \( 1, 2, 3, \ldots \) the symbols of \( \mathcal{A} \) and in bold red \( 1, 2, 3, \ldots \) the symbols of \( \mathcal{B} \) ordered with \( 1 < 2 < 3 < \cdots < 1 < 2 < 3 < \cdots \). Then,

\[
G_{132}(\mathcal{A} \sqcup \mathcal{B}) = 121 + 131 + 132 + \cdots + 111 + 112 + \cdots + 121 + 131 + \cdots + 121 + \cdots
\]
\[
= 121 + 131 + 132 + \cdots + 11 \cdot 1 + 11 \cdot 2 + 12 \cdot 1 + \cdots + 1 \cdot 21 + 1 \cdot 31 + \cdots + 123 + 132 + 232 + \cdots
\]
\[
= \Delta(G_{132}) = 1 \otimes G_{132} + G_1 \otimes G_{21} + G_{12} \otimes G_1 + G_{132} \otimes 1.
\]

3.3 Good Hopf algebras

We call a Hopf algebra \( \mathcal{H}_\varphi \) good if it is defined by a Hopf polynomial realization \( r_\varphi \). We call a function \( \varphi \) good if it produces a good Hopf algebra \( \mathcal{H}_\varphi \). Currently, we know three main good Hopf algebras: FQSym, WQSym and PQSym are respectively associated to the standardization, packing and parkization functions.

4 Good monoids

In the previous section (Section 3), we realized some Hopf algebras in free algebras. In this section, we give sufficient conditions on a congruence \( \equiv \) to build a combinatorial quotient of a good Hopf algebra. We call a monoid good if it satisfies these conditions. The first condition is about the \( \varphi \)-map used to realize the good Hopf algebra in free algebras. We give a sufficient compatibility between \( \varphi \) and \( \equiv \) to ensure the product is carried to the quotient. The second condition ensures that the alphabet doubling trick map. It is used to project the coproduct in the quotient. Under these conditions, a monoids is guaranteed to produce a Hopf algebra quotient (Theorem 1). Furthermore, these conditions on monoid are preserved under taking infimum and supremum (Theorem 2).

4.1 Definition

The notion of Good monoids has been introduced by Hivert-Nzeutchap [Hiv04, HN07] to build quotients (sub-algebras) of FQSym. We could also mention PhD thesis [Gir11b].
A good monoid is a monoid which has similar properties, as the plactic monoid [LS81, Knu73]. We consider a free monoid \( \mathfrak{A}^* \) with concatenation product \( \cdot \), a congruence \( \equiv \) on \( \mathfrak{A}^* \) and a map \( \varphi : \mathfrak{A}^* \rightarrow \mathfrak{A}^* \). We define the evaluation \( ev(w) \) of a word \( w \) as its number of occurrences of each letter of \( w \). For example, the words \( ejajv \) and \( jjaev \) have the same evaluation: both have one \( a \), one \( e \), one \( v \) and two \( j \). The free monoid \( \mathfrak{A}^*/\equiv \) is a \( \varphi \)-good monoid if it has the following properties:

**Definition 3 (\( \varphi \)-congruence):** The congruence \( \equiv \) is a \( \varphi \)-congruence if for all \( u, v \in \mathfrak{A}^* \), \( u \equiv v \) if and only if \( \varphi(u) \equiv \varphi(v) \) and \( ev(u) = ev(v) \).

This first compatibility is sufficient to build a quotient algebra of \( \mathfrak{A}_\varphi \).

**Definition 4 (Compatibility with restriction to alphabet intervals):** The congruence \( \equiv \) is compatible with the restriction to alphabet intervals if, for all \( u, v \in \mathfrak{A}^* \) such that \( u \equiv v \), one has \( u_I \equiv v_I \) for any \( I \) interval of \( \mathfrak{A} \), where \( w_I \) is word restricted to the alphabet \( A \).

This second compatibility in association with the first ensures that alphabet doubling trick defines a quotient coproduct. Both compatibilities give us an extended definition of a Hivert-Nzeutchap’s good monoid which one is defined only with \( \varphi \) the standardization map:

**Definition 5 (\( \varphi \)-good monoid):** A quotient \( \mathfrak{A}^*/\equiv \) of the free monoid is a \( \varphi \)-good monoid if \( \equiv \) is a \( \varphi \)-congruence and is compatible with restriction to alphabet intervals. We call such a congruence a \( \varphi \)-good congruence.

In the following examples, we denote words of \( \mathfrak{A}^* \) by \( u, v, w \) and the letters by \( a, b, c \).

**Example 3 (sylvester and stalactic monoids):** The sylvester congruence: \( \equiv_{sylv} \), defined by

\[
    u \cdot ac \cdot w \cdot b \cdot v \equiv_{sylv} u \cdot ca \cdot w \cdot b \cdot v \quad \text{whenever } a \leq b < c ,
\]

is \( std \)-compatible and compatible with the restriction to alphabet intervals. Thanks to the binary search tree insertion algorithm the equivalence classes are in natural bijection with binary search trees. The quotient monoid is a monoid on binary search trees called the sylvester monoid in [HNT04, HNT05].

The stalactic congruence [HNT08a]: \( \equiv_{stal} \), defined by

\[
    u \cdot ba \cdot v \cdot b \cdot w \equiv_{stal} u \cdot ab \cdot v \cdot b \cdot w ,
\]

is compatible with packing but not with standardization. The quotient monoid is the stalactic monoid. It is clear that any stalactic class contains a word of the form \( \cdots a_i^m \cdots a_i^{m_k} \), where the \( a_i \) are distinct. We call these words canonical. We represent a stalactic class with a planar diagram such that, in any column, the boxes contain the same letter.

\[
\begin{array}{cccc}
3 & 1 & 2 & 1 \downarrow \\
5 & 1 & 1 & 1 \downarrow \\
\end{array}
\]

\[
51543151145312455 \equiv_{stal} 3^21^521^43^56 \leftrightarrow \begin{array}{cccc}
3 & 1 & 2 & 1 \downarrow \\
5 & 1 & 1 & 1 \downarrow \\
\end{array}
\]

4.2 Hopf algebra quotient

These different good monoids tools was used to (re-)define several Hopf algebra quotients: \( \text{FSym} \) the Free Symmetric functions Hopf algebra [DHT02], \( \text{PBT} \) [LR98, HNT04, HNT05] or Baxter Hopf algebra [Gir11a, Gir12]; the Hopf algebra associated with the stalactic monoid [HNT08a]; or \( \text{CQSym} \) [NT04, NT07] (a \( \text{PQSym} \) quotient).
Lemma 1 (Algebra quotient): Let $H_\varphi$ be a good Hopf algebra and $\equiv$ be a $\varphi$-good congruence such that its free monoid quotient is a $\varphi$-good monoid. Then, the quotient $H_\varphi/\equiv$ is an algebra quotient whose bases are indexed by $\text{can}_{\varphi}/\equiv$, identifying basis elements $m_u$ and $m_v$ whenever $u \equiv v$.

Example 4 (PBT and Hopf algebra stalactic): We go back to Example 3. The sylvester quotient of $\text{FQSym}$ is the Hopf algebra $\text{PBT}$ [LR98, HNT04, HNT05].

The stalactic monoid gives a quotient of $\text{WQSym}$. Let $\pi$ be the projection of $\text{WQSym}$ in $\text{WQSym}/\equiv_{\text{stal}}$ and $u := 112$ and $v := 11$ two (packed) words. We denote by $\pi$ the projection of $M_u$ by $Q_s$, with $s$ the planar diagram associated to the stalactic class of $u$.

\[
\pi(M_{112} \times M_{11}) = Q_{12} \times Q_{11} = Q_{12} + Q_{11} + Q_{12} + Q_{11} + Q_{12} + Q_{11} + Q_{12} + Q_{11}
\]

Lemma 2 (Coalgebra quotient): The quotient $H/\equiv$ is a coalgebra quotient.

Sketch of the Proof: The relation $\equiv$ is compatible with the restriction to alphabet intervals, hence the alphabet doubling trick ensures that coproduct projects to the quotient.

Example 5:

\[
\pi(\Delta(M_{332122})) = \pi(1 \otimes M_{332122} + M_1 \otimes M_{22111} + M_{2122} \otimes M_{11} + M_{332122} \otimes 1) = \Delta(Q_{12} \otimes Q_{11} + Q_{11} \otimes Q_{12} + Q_{12} \otimes Q_{11} + Q_{11} \otimes Q_{12} + Q_{12} \otimes 1)
\]

Theorem 1 (Good monoid and good Hopf algebra): Let $H_\varphi$ be a good Hopf algebra and $\equiv$ be a $\varphi$-good congruence. The quotient $H/\equiv$ is a Hopf algebra quotient.

Corollary 1: The dual Hopf algebra $(H/\equiv)^\#$ is a sub-algebra of the dual Hopf algebra $H^\#$, with basis given by:

\[
\sum_{u \in U} m_u^\# = \sum_{u \in U} m_u^\# .
\]

4.3 Operations

Previously we introduced some good functions $\varphi$: $\text{std}$, $\text{pack}$ (and $\text{park}$). It is interesting to investigate the connections between them:

Definition 6 (refinement): Let $\varphi$ and $\pi$ be two functions. We say that $\pi$ refines $\varphi$, written $\varphi < \pi$ if $\varphi(\pi(u)) = \varphi(u)$ for all $u \in \mathbb{A}^*$.

It is clear that refinement is an order.

Proposition 2 ($\text{std}$, $\text{tass}$, $\text{park}$ and refinement): For these three functions: standardization $\text{std}$, packing $\text{pack}$ and parking $\text{park}$ we have the relation: $\text{std} \prec \text{pack} \prec \text{park}$.

Proposition 3 (Good functions and refinement): Let $\varphi$ and $\pi$ be two good functions such that $\varphi < \pi$. Then any $\varphi$-good monoid is a $\pi$-good monoid.
Figure 1: We start by considering the packed word $45142234212$, and insert it in a BSTM by the algorithm $P$; that gives us $P(45142234212)$ above in the middle. On the right, there is the BTm ($B_m(w)$) associated with the BSTM $P(w)$ of $WQSym/\equiv_t$. At the top of the figure there is the $P$-symbol given by $P$ or $B_m$ and below the $Q$-symbol is given by $Q$.

Propositions 2 and 3 give us, for example, that any $std$-good monoid is $pack$-good. Furthermore operations on two $good$ congruences give $good$ congruences.

**Theorem 2 ($\lor$, $\land$ and $good$ congruences):** The union and intersection of two $\varphi$-good congruences $\sim$ and $\approx$ are $\varphi$-good congruences.

As an intriguing consequence the lattice structure on monoids is transported to Hopf algebras. Several examples of this are known.

**Example 6:** The intersection ($\equiv_{sylv} \land \equiv_{#sylv}$) of the sylvester relation (5) and its image under the SCHÜTZENBERGER involution gives $std$-good monoid: the Baxter monoid [Gir11a, Gir11b].

The union ($\equiv_{sylv} \lor \equiv_{#sylv}$) of those relations gives the hypoplactic monoid [Nov00].

In the sequel, we study in detail another example.

5 The union of the sylvester and the stalactic congruences

As an application of the preceding construction, we consider the union ($\equiv_{sylv} \lor \equiv_{stal}$) of the sylvester congruence (5) and the stalactic congruence (6); we call it the taïga relation $\equiv_t$.

\[
\begin{align*}
    u \cdot ac \cdot v \cdot b \cdot w & \equiv_t u \cdot ca \cdot v \cdot b \cdot w \quad \text{for } a \leq b < c, \\
    u \cdot ba \cdot v \cdot b \cdot w & \equiv_t u \cdot ab \cdot v \cdot b \cdot w
\end{align*}
\]

From Proposition 3 we know that the sylvester congruence (5) is a $pack$-good congruence and from Theorem 2 we deduce that the taïga monoid is a $pack$-good monoid.

5.1 Algorithm and taïga monoid

The taïga congruence can be calculated using an insertion algorithm similar to the binary search tree insertion (see Algorithm 3 for a definition). This insertion algorithm uses a search tree structure:

**Definition 7 (Binary search tree with multiplicity):** A (planar) binary search tree with multiplicity (BS-TM) is a binary tree $T$ where each node is labelled by a letter $l$ and a non-negative integer $k$, called the multiplicity, so that $T$ is a binary search tree if we drop the multiplicities and such that each letter appears at most once in $T$.

We denote by $(l, k)$ a node label and for any node $n$, by $l(n)$ its letter and by $m(n)$ its multiplicity.
Algorithm 3: insertion in a BSTM

Data: $t$ a BSTM with $L_t$ and $R_t$ its left and right subtrees, and $l$ a letter of $A$

Result: $t$ with $l$ inserted

1. if $t$ is empty tree then
2. $t \leftarrow$ node labelled by $(l, 1)$
3. else
4. if $l(t) = l$ then
5. increment $m(t)$
6. else
7. if $l(t) < l$ then insert recursively $l$ in $L_t$
8. else insert recursively $l$ in $R_t$
9. return $t$

We denote by $P(w)$ the result of the insertion using Algorithm 3 of $w$ from the right to the left in the empty tree (cf. the left part of the figure).

Proposition 4: The taiga classes are the fibers of $P$. That is for $u$ and $v$ two words: $u \equiv_t v$ if and only if $P(u) = P(v)$.

The $Q$-symbol of $w$ is the tree $Q(w)$ of same shape as $P(w)$ which records the positions of each inserted letter. This gives us a Robinson-Schensted like correspondence [Lot03] (cf. Figure 1). As a corollary of Theorem 2 we get

Corollary 2: The taiga monoid is a tass-good monoid.

5.2 Quotient of $WQSym$: $PBTm$

Thanks to Proposition 4, the set of packed words giving the same tree by algorithm $P$ is exactly a taiga class of packed words. As in [HNT04], we consider a binary trees with multiplicities without letters.

Definition 8 (BTM): A binary tree with multiplicities (BTM) is a (planar) binary tree labelled by non-negative integers on its nodes. The size of a BTM $T$ denoted by $|T|$ is the sum of the multiplicities.

Let $T_w$ be a BSTM associated to a packed word $w$, and $T$ be the BTM obtained by removing its letters. One can recover uniquely $T_w$ from $T$: indeed each letter of $T_w$ is deduced by a left infix reading of $T$. We identify the set of words in $\text{pack}(A^\ast)/\equiv_t$ of size $k$ (for $k \geq 0$) with the set of BTM of size $k$. We denote $B_m$ the algorithm which computes the BTM associated to the BSTM computed by $P$ (cf. Figure 1).

Let us denote by $S(t)$ the generating series of these trees counted by size. The generating serie satisfies the following functional equation (see A002212 of Slo):

$$ S(t) = 1 + \frac{S(t)^2}{1 - t} \quad \text{and thus} \quad S(t) = \frac{1 - t - \sqrt{5t^2 - 6t + 1}}{2t} $$

$$ = 1 + t + 3t^2 + 10t^3 + 36t^4 + 137t^5 + 543t^6 + 2219t^7 + \ldots $$
This structure is in bijection with binary unary tree structure. Here is the list of trees of size 0,1,2 and 3:

```
        1
   2-1-3-1
         4
```

```
        1
   3-1-2-1
         4
```

```
        1
   2-1-3-1
         4
```

```
```

With Lemma \[\text{I}\] and Theorem \[\text{I}\] we know that the quotient of \text{WQSym}(\mathcal{R}) by the ta\"iga relations has a natural basis indexed by \text{taass}(\mathcal{R}^*)/\Xi, identified by BTM. We call \text{PBTm} (planar binary tree with multiplicities) that quotient. More precisely, we consider the basis \((\mathcal{M}_u)_u\) of \text{WQSym} obtained by the Hopf polynomial realization \(r_{\text{taass}}\). We denote by \((Q^m_t)_t\) the canonical projection by the map \(\pi\) of \((\mathcal{M}_u)_u\) in \text{PBTm} such that \(\pi(\mathcal{M}_u) := Q^m_t\) if \(t = B_m(u)\). The product and coproduct are given by some explicit algorithms. For brevity, we only give here some examples.

The product on the \((Q^m_t)_t\) basis is described thanks to the projection \(\pi\). For example,

\[\pi(\mathcal{M}_{1312} \times \mathcal{M}_1) = \pi(\mathcal{M}_{13121} + \mathcal{M}_{13122} + \mathcal{M}_{13123} + \mathcal{M}_{13124} + \mathcal{M}_{14123} + \mathcal{M}_{14132} + \mathcal{M}_{2423}) = Q^m_m \times Q^m_m = Q^m_m + Q^m_m + Q^m_m + Q^m_m + Q^m_m + Q^m_m + Q^m_m .\]

Similarly, the coproduct is described thanks to the projection \(\pi\). We could remark that it is the same as the coproduct on \((Q^m_t)_t\) of \text{PBT} keeping the multiplicities with nodes. For example,

\[\pi(\Delta(\mathcal{M}_{1312})) = \pi(1 \otimes \mathcal{M}_{1312} + \mathcal{M}_{11} \otimes \mathcal{M}_{21} + \mathcal{M}_{112} \otimes \mathcal{M}_{1} + \mathcal{M}_{3112} \otimes 1) = \Delta(Q^m_m) = 1 \otimes Q^m_m + Q^m_m \otimes Q^m_m + Q^m_m \otimes Q^m_m + Q^m_m \otimes Q^m_m \otimes 1 .\]

We consider \text{PBTm^\#} the dual of \text{PBTm}. This is a sub-algebra of \text{WQSym^\#}. We denote by \((P^m_t)_t := (Q^m_t)^\#\) its dual basis: \((Q^m_t, P^m_t) = \delta_{t,t}^\circ\). The product is given by:

\[P^m_t \times P^m_t'' = \sum_t (\Delta(Q^m_t), P^m_{t''} \times P^m_t) P^m_t'' . \tag{10}\]

Here is an example,

\[
P^m_m \times P^m_m = P^m_m + P^m_m + P^m_m + P^m_m + P^m_m + P^m_m .
\]

If we consider only shape tree, the product is exactly the product of \((P_t)_t\) basis in \text{PBT} \cite{HNT04}. Hence this product is a shifted shuffle on trees.

The coproduct is given by: \(\Delta^\#(P^m_t) = \sum_{t',t''} (Q^m_t \times Q^m_{t''}, P^m_t) P^m_{t'} \otimes P^m_{t''}\). Here is an example:

\[
\Delta^\#(P^m_m) = 1 \otimes P^m_m + P^m_m \otimes \left(P^m_m + P^m_m\right) + \left(P^m_m + P^m_m\right) \otimes P^m_m + P^m_m \otimes 1 .
\]
6 The hook length formula

It is well known from [Knu73] (§5.14 ex. 20) that the number of decreasing labelling of a binary tree is given by a simple product formula. [HNT04] remarks that this is also the number of permutations given upon a tree by the binary search tree insertion. In this section we generalize this formula for trees with multiplicities.

**Proposition 5:** The cardinal \( f(T) \) of the taïga class associated to \( T \) (i.e. the set of packed words giving the tree \( T \) by the insertion algorithm \( \mathcal{B}_m \)) is given by

\[
f(T) = |T|! \left( \prod_{t \in T} |t| (m(t) - 1)! \right)^{-1}.
\]

(11)

where \( t \) ranges through all the subtrees of \( T \) and \(|T|\) denotes the size of \( T \) (the sum of the multiplicities).

**Example 7:** The taïga class of \( T := \begin{array}{cc}
1 & 2 \\
\end{array} \begin{array}{c}
2 \\
\end{array} \begin{array}{c}
2 \\
\end{array} \) contains 12 packed words \( w \):

\[
23132, 33122, 31232, 32312, 13232, 23312, 32132, 21332, 31322, 12332, 13322.
\]

The class of \( \begin{array}{c}
2 \\
\end{array} \begin{array}{cc}
2 & 4 \\
\end{array} \begin{array}{c}
3 \\
\end{array} \begin{array}{c}
2 \\
\end{array} \) contains 23,377,600 = \( \frac{18!}{(18 \cdot 9 \cdot 7 \cdot 7 \cdot 4 \cdot 2)(1!1!6!0!3!1!)} \) packed words.

This formula is easily proven by induction. However, we prefer to give a generating series proof as in [HNT08b]. Let \( \mathcal{A} \) be an associative algebra, and consider the functional equation for power series \( x \in \mathcal{A}[[z]] \):

\[
x = a + \sum_{k \geq 1} B_k(x, x),
\]

(12)

where \( a \in \mathcal{A} \) and for any \( k > 0 \), \( B_k(x, y) \) is a bilinear map with values in \( \mathcal{A}[[z]] \). We suppose such that the valuation of \( B_k(x, y) \) is strictly greater than the sum of the valuations of \( x \) and \( y \) (plus \( k \)). Then Equation (12) has a unique solution:

\[
x = a + \sum_{k \geq 1} \left( B_k(a, a) + B_k(a, \sum_{k' \geq 1} B_{k'}(a, a)) + B_k(\sum_{k' \geq 1} B_{k'}(a, a)), a) + \ldots \right)
\]

(13)

where for a tree \( T \), \( B_T(a) \) is the result of evaluating the expression formed by labelling by \( a \) the leaves of the complete tree associated to \( T \) and by \( B_k \) its internal node labelled by \( k \).

For example: \( B \begin{array}{cc}
2 & 2 \\
\end{array} \begin{array}{c}
2 \\
\end{array} \begin{array}{c}
6 \\
\end{array} \begin{array}{c}
2 \\
\end{array} \) = \( B_3(B_6(a, B_2(a)), B_2(a, a)) \).
So if we try to solve the fixed point problem:

\[ x = 1 + \int_0^z e^s x(s)^2 ds = 1 + \sum_{k \geq 1} \int_0^z \frac{s^{k-1}}{(k-1)!} x(s)^2 ds = 1 + \sum_{k \geq 1} B_k(x, x), \tag{14} \]

where \( B_k(x, y) = \int_0^z \frac{s^{k-1}}{(k-1)!} x(s) y(s) ds \). Then for a binary tree of non-negative integer \( T \), \( B_T(1) \) is the monomial obtained by putting 1 on each leaf and integrating at each node \( n \) the product of the evaluations of its subtrees and \( s^k/k! \) with \( m(n) = k + 1 \).

For example:

\[ \begin{array}{c}
\text{1} \\
\text{2} \\
\text{3} \\
\text{1} \\
\end{array} \rightarrow \begin{array}{c}
\text{z} \\
\text{1} \\
\text{1} \\
\text{1} \\
\text{1} \\
\end{array} \]

One can observe that \( B_T(1) = f(T) \frac{z^n}{n!} \), where \( n = |T| \).

To prove the hook length formula, following the same technique as in \[HNT08\], we want to lift in \( \text{WQSym}^\# \) the fixed point computation of Equation 14. Recall from \[Hiv99, NT+06\] that the multiplication rule of the dual basis \( \mathbb{S}_u (\mathbb{M}^\#_u := \mathbb{S}_u) \) is, for \( u, v \) two packed words of respectively size \( k \) and \( l \),

\[ \mathbb{S}_u \cdot \mathbb{S}_v = \sum_{w \in (u) \sqcup (v)} \mathbb{S}_w \] with \( u \sqcup v \) the set of words appearing in the shifted shuffle of \( u \) and \( v \). The number of terms is a binomial: \( \binom{k+l}{k} \) for \( k \) and \( l \) the length of \( u \) and \( v \). Hence, the linear map \( \phi : \mathbb{S}_u \mapsto \frac{z^n}{n!} \) with \( n \) the length of \( u \) is a morphism of algebras from \( \text{WQSym}^\# \) to \( \mathbb{K}[[z]] \). For \( u, v \) two packed words of respective size \( n - 1 \) and \( m \), set

\[ B_k(\mathbb{S}_u, \mathbb{S}_v) := \sum_{w \in (u) \sqcup (1 \cdots k - 1 \sqcup v) \sqcup n} \mathbb{S}_w. \tag{15} \]

The crucial observation which allows to express the hook length formula in a generating series way is the following theorem:

**Theorem 3:** In the binary tree (with multiplicity) solution (Equation 13) of Equation 14,

\[ B_T(1) = \sum_{B_m(u) = T} \mathbb{S}_u, \tag{16} \]

In particular, \( B_T(1) \) coincide with \( \mathbb{P}^T \), the natural basis of \( \text{PBTm}^\# \).

**Corollary 3:** The number of packed words \( u \) such that \( B_m(u) = T \) is computed by \( f(T) \).

## 7 Conclusion, work in progress and perspectives

In this paper, we unraveled some new combinatorics on binary trees with multiplicities from the union of the sylvester and stalactic monoids. Using the machinery of realizations, we built a Hopf algebra on those trees, allowing us to give a generating series proof of a new hook length formula. Following \[HNT08\], it is very likely that we will also be able to prove a \( q \)-hook length formula. On the other hand, the usual case of the LODAY-RONCO algebra has a lot of nice properties. For example, the product and coproduct can be expressed by the means of an order on the trees called the Tamari Lattice \[LR98\]. It would be good to
know if such a lattice exists for trees with multiplicities. This should also relate to N. READING work on lattice congruences [Rea05]. Also it could be interesting to study some other combinations in the lattice of good monoids. For example, the union of the plactic monoid and the stalactic monoid should give a Hopf algebra of tableaux with multiplicities. Finally, in our construction, it seems that \( \text{std}, \text{tass} \) and \( \text{park} \) play some canonical role from which everything else is built. Are there some more examples? Is there a definition for such a \( \varphi \)-map? Could we except to always have a hook formula as soon as we have a good monoid?

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