On Ambiguity in Knot Polynomials for Virtual Knots

A.Morozov\textsuperscript{a,b,c}, And.Morozov\textsuperscript{a,b,c,d} and A.Popolitov\textsuperscript{a,b,e}

\textsuperscript{a} ITEP, Moscow 117218, Russia
\textsuperscript{b} Institute for Information Transmission Problems, Moscow 127994, Russia
\textsuperscript{c} National Research Nuclear University MEPhI, Moscow 115409, Russia
\textsuperscript{d} Laboratory of Quantum Topology, Chelyabinsk State University, Chelyabinsk 454001, Russia
\textsuperscript{e} Korteweg-de Vries Institute for Mathematics, University of Amsterdam, P.O. Box 94248, 1090 GE Amsterdam, The Netherlands

ABSTRACT

We claim that HOMFLY polynomials for virtual knots, defined with the help of the matrix-model recursion relations, contain more parameters, than just the usual $q$ and $A = q^N$. These parameters preserve topological invariance and do not show up in the case of ordinary (non-virtual) knots and links. They are most conveniently observed in the hypercube formalism: then they substitute $q$-dimensions of certain fat graphs, which are not constrained by recursion and can be chosen arbitrarily. The number of these new topological invariants seems to grow fast with the number of non-virtual crossings: 0, 1, 1, 5, 15, 91, 784, 9160, ... This number can be decreased by imposing the factorization requirement for composites, in addition to topological invariance – still freedom remains. None of these new parameters, however, appear in HOMFLY for Kishino unknot, which thus remains unseparated from the ordinary unknots even by this enriched set of knot invariants.

1 Introduction

The most efficient methods to calculate knot/link polynomials \cite{1,2} are actually dealing with link diagrams – oriented graphs of valence $(2,2)$ with two types (colors) of vertices. In the Chern-Simons theory \cite{3} this description appears in the temporal gauge $A_0 = 0$ \cite{4} (while a choice of holomorphic gauge $A_1 + iA_2 = 0$ leads to the formalism of Kontsevich integrals \cite{5}). Topological invariance is then provided by the Reidemeister invariance. For ordinary knots and links graphs are planar, non-planar graphs are interpreted as associated with virtual knots/links, provided the set of Reidemeister moves is appropriately enlarged \cite{6}. Planar graphs are intimately related to braids, what allows to apply the powerful theory of Hecke algebras – which is in the base of the modern (universal) versions \cite{7,8} of the Reshetikhin-Turaev (RT) formalism \cite{9}. This is not the case for non-planar graphs, and these methods are not directly applicable (though available answers imply that just a minor – but still unknown – modification can be needed). Therefore virtual knot/link polynomials need to be studied by different methods. Historically first was the artful skein relation method \cite{6} (which is normally a small part of the RT formalism, restricted to the fundamental representation \cite{7}). A more systematic alternative approach is the hypercube method of \cite{10,11}, which was successfully applied to virtual knots in \cite{12} and \cite{13}. Moreover, in \cite{13} it was reformulated in matrix model terms, what implies applicability of topological recursion ideas \cite{14} – this seems to be the way, explaining emergence of skein relations in simple situations from the hypercube formalism. In fact the relations, found in \cite{13} (which are probably equivalent to the older MOY relations \cite{15}), are not quite recursive (what is also the case for the skein relations) – still, in combination with enumerative methods, they are sufficient to calculate fundamental HOMFLY for ordinary knots (this is currently checked up to 10 intersections – for the whole Rolfsen table). However, application of the same computer program (available at \cite{16}) to virtual knots does not fix polynomials unambiguously. It is the purpose of this paper to describe this (empirical) situation in some detail.

In the formalism of \cite{13} the basic objects are “quantum dimensions”, associated with fat graphs, and above-mentioned relations are associated with graph reshufflings. What happens is that dimensions for some of the
graphs remain undetermined – but this never happens to graphs, appearing in consideration of ordinary (non-virtual) knots and links. A natural hypothesis is that we actually encounter new topological invariants, which are not observable in the world of ordinary knots. We did some checks ensuring that the new parameters are indeed topological invariants – but no proof is yet available.

One could suspect that such additional invariants exist – already because of the old puzzle of Kishino virtual knot [17, 6], which is a composite of two unknots, not reducible to unknot by the Reidemeister-Kauffman moves. While Kishino knot polynomial does not depend on any of our new parameters – and thus remains undistinguished from unknot by fundamental HOMFLY – new parameters do appear for a closely related knot. This knot has the same shape of a planar diagram but all intersections black, so it is a composite of two virtual trefoils. Not only does its HOMFLY contain new parameters, it also contains more new parameters, then the virtual trefoil itself.

The paper is organized as follows. In Section 2 we remind how to calculate HOMFLY polynomial with help of skein relations. In Section 3 we show, that if we naively apply this method to virtual knots, we are faced with ambiguities. Then in Sections 4 and 5 we develop a proper language to address these issues – the hypercube formalism. In Section 6 we restate the ambiguity in this language, where it takes the form of new parameters – free values of some quantum dimensions. In Section 7 we describe an algorithm, that can check, to what extent keeping these parameters free is consistent with topological invariance. In Section 8 we present the outcome of this algorithm’s work, so impatient reader may immediately look here. In Section 9 we study, whether the ambiguity can be constrained by imposing factorization of HOMFLY polynomial for composite knots. We finish the main body of the paper by formulating two conjectures about these new empirically observed free parameters. In Appendix A we consider in detail example of Kishino knot – showing that new parameters break factorization of composites. Appendix B contains more statistics on the results obtained by algorithm of section 7. Finally, in Appendix C we show an alternative line of development of the hypercube formalism – using symmetric representation [2] instead of antisymmetric [1, 1].

2 Skein method = skein relations + Reidemeister moves

One of the ways to calculate fundamental HOMFLY polynomial for a given planar diagram of a knot is to use the skein relation

\[ A \biggl( \begin{array}{c} \biggr) - A^{-1} \biggl( \begin{array}{c} \biggr) = z \biggl( \begin{array}{c} \biggr) \biggr) \biggr) \] (1)

where \( A = q^N \) and \( z = q - q^{-1} \). From the point of view of RT approach this relation is the property of quantum \( R \)-matrix in the fundamental representation

\[ \left( R - \frac{1}{Aq} \right) \left( R + \frac{q}{A} \right) = 0 \] (2)

(for higher representations \( R \)-matrix has more different eigenvalues, and the story gets more involved).

The method is to gradually express HOMFLY polynomial for a diagram through HOMFLY polynomials of simpler diagrams, eventually expressing everything through completely unknotted diagrams.

Each simplification step consists of two substeps:

- skein relation is applied to some crossings of a planar diagram
- some of resulting summands are simplified with the help of Reidemeister moves

Crucial is the second substep, without it the procedure would not terminate at all, because one of the summands after the application of relation always has the same number of crossings as in the original diagram. But can we always simplify this summand with help of Reidemeister moves, such that it contains less crossings, than the diagram we started with?

The answer is yes – if we choose crossings, at which skein relations are applied, in a clever way. Namely, for non-virtual knots it is easy to observe that

**Proposition 2.1.** Every planar diagram can be unknotted by changing some of its crossings for inverse ones.

Thus, if we choose to apply skein relation precisely at all these “unknotting” crossings, then the problematic summand with non-decreased number of crossings can definitely be simplified – even completely unknotted!
with the help of Reidemeister moves – because it is the unknot. Thus in principle the procedure can calculate arbitrary fundamental HOMFLY.

The problem in practice is to recognize, which crossings are the unknotting ones? If we are not concerned with efficiency and just want to show that algorithm terminates in finite time, we can choose the most naive, brute-force enumerative approach.

We just loop over all $2^{#crossings}$ possible choices and check, whether diagram we obtain by changing these crossings is an unknot. To check, whether some diagram is an unknot is a famous “unknotting problem”. Again we can choose the most naive brute-force approach to it.

The number of Reidemeister moves, needed to unknot any planar diagram, is at most exponential in the number of crossings \[18\]. Hence, the full search over all possible sequences of Reidemeister moves (of restricted length) will take at most doubly exponential time – but finite nonetheless. Hence, skein method will give an answer for any planar diagram in finite time as well.

Of course, in practice, it is possible to optimize this brute-force procedure in many ways. First of all, it is very often possible to apply skein relation at just one crossing, but in such a way, that the number of crossings of the problematic summand can be immediately reduced with help of the second Reidemeister move. Second, there is a lot of smarter approaches to the unknotting problem, which are at most exponential in the number of crossings (though it is still unknown, whether there is a polynomial algorithm for that).

**Example: trefoil** Let’s apply the skein method to get HOMFLY polynomial of a trefoil

\[
\begin{array}{c}
\includegraphics{trefoil.png}
\end{array}
\]

First we substitute the upper black vertex by the white one, using skein relation, to get

\[
\begin{array}{c}
\includegraphics{trefoil_2.png}
\end{array}
= A^2 \begin{array}{c}
\includegraphics{trefoil_3.png}
\end{array} - Az \begin{array}{c}
\includegraphics{trefoil_4.png}
\end{array}
\]

We see, that the diagram in the first item still has the same number of crossings – 3. But we can apply the sequence of 2nd and 1st Reidemeister moves to unknot it – thus it equals $[N] = q^2 - q^{-2} = \frac{A - A^{-1}}{q - q^{-1}}$, the quantum dimension of the fundamental representation, which is usually associated with the unknot (this convention defines “non-normalized” HOMFLY, if one instead associates just unity with the unknot, then all polynomials should be divided by $[N]$ to provide “normalized” HOMFLY). The diagram in the second summand is a Hopf link and has 2 crossings, so it’s simpler, than the original one – we can say, that we’ve calculated it on the previous step of the recursion, and it’s equal to $A^2[N]^2 - Az[N]$. So the full answer is

\[
H_{\text{trefoil}} = A^2[N] - Az \left(A^2[N]^2 - Az[N]\right) = (-A^4 + A^2z^2 + 2A^2) \cdot [N]
\]

Though “unknotting” proposition 2.1 is trivial for usual knots, this statement often fails for virtual knots, as we shall shortly see in the next section. Therefore, if we apply skein method to virtual knots, we can not always get a number (i.e. a Laurent polynomial in $A$ and $q$) – sometimes some un-simplifiable planar diagrams remain in the answer. This, however, is not a bad thing: from a proper point of view (hypercube formalism) these un-reducible pictures are free parameters, that do not conflict with topological invariance. Hence, they can be treated as new parameters and HOMFLY polynomial just starts to depend on a larger (very likely, infinite) set of variables – and still remains topological invariant.

### 3 Skein method does not completely determine HOMFLY for virtual knots: example of Hopf link

Let us now apply the skein method to the simplest virtual planar diagram – the virtual Hopf link

\[
\begin{array}{c}
\includegraphics{hopf.png}
\end{array}
\]

Hence, skein method will give an answer for any planar diagram in finite time as well.
We immediately get into a trouble: the skein relation

\[ A^{-1} \begin{array}{c} \circ \hspace{0.5cm} \circ \end{array} = A \begin{array}{c} \circ \hspace{0.5cm} \circ \end{array} - z \begin{array}{c} \circ \hspace{0.5cm} \circ \end{array} \] (5)

does not express HOMFLY for virtual Hopf through HOMFLY for simpler knot \( s! \) Instead one of the items at the right hand side is mirror reflection of original one at the l.h.s. Even assuming that they are related by the symmetry \( q \rightarrow 1/q \) (that sends \( A \rightarrow A^{-1} \) and \( z \rightarrow -z \)), we conclude, that the most what we can do is to express part of HOMFLY for virtual Hopf as a number

\[ A^{-1} \begin{array}{c} \circ \hspace{0.5cm} \circ \end{array} q^{-1} - A \begin{array}{c} \circ \hspace{0.5cm} \circ \end{array} = z[N] \] (6)

while another linear combination

\[ A^{-1} \begin{array}{c} \circ \hspace{0.5cm} \circ \end{array} q + A \begin{array}{c} \circ \hspace{0.5cm} \circ \end{array} \frac{1}{q} \]

remains unconstrained by skein relation.

Of course, one may wonder, whether it is possible to first twist the planar diagram of a virtual Hopf link in some way, and only then apply skein method, such that in the end we would still get a number. This question is surprisingly hard to answer and is, in a sense, the central question of this paper. The best we can do now is to present checks, that such a twisting is not possible, when we consider diagrams with up to and including 7 crossings. To address the issue systematically, we switch in the next section to the hypercube formalism.

4 Hypercube formula: solution to skein relation

In [13] we wrote the hypercube-style formula, that expressed HOMFLY polynomial for a virtual knot as a sum over amputations of edges of its \( \text{dessin d’enfant} \). \( \text{H}_{\mathcal{L}} \)

\[ H_{\mathcal{L}} = q^{(N-1)(n_\bullet(\Gamma_{\mathcal{L}})-n_\circ(\Gamma_{\mathcal{L}}))} \sum_{\gamma \subseteq \Gamma_{\mathcal{L}}} (-q)^{n_\bullet(\gamma)-n_\circ(\gamma)} \cdot D_\gamma(q, N) \] (7)

Here \( \Gamma_{\mathcal{L}} \) is the fat graph, that is associated to the planar diagram \( \mathcal{L} \). Vertices of this fat graph are Seifert cycles of a planar diagram and edges are its crossings (in [13] everything is explained in more detail, with examples). \( n_\bullet \) and \( n_\circ \) are number of black and white edges, respectively. Quantum dimensions \( D_\gamma \) are functions of \( q \) and \( N \), that are associated to any fat graph. The sum in the formula runs over all subgraphs \( \gamma \) of \( \Gamma_{\mathcal{L}} \), such that they contain all the vertices of \( \Gamma_{\mathcal{L}} \), but may not contain some edges (hence the name “sum over amputations”).

In fact, this formula almost appears already in the old work [15] – but for planar fat graphs only. In this case the formula can be easily explained from RT point of view. Namely, if one takes expressions for fundamental \( R \)-matrix and its inverse through the projector onto antisymmetric representation \( P_{[1,1]} \)

\[ R = q^{-N} \cdot \left( \frac{1}{q} \cdot I \otimes I - [2]P_{[1,1]} \right) \]

\[ R^{-1} = q^N \cdot (q \cdot I \otimes I - [2]P_{[1,1]}) \] (8)

then the hypercube formula [7], is just the result of expanding all the brackets in the RT tensor contraction. The \( q \)-dimensions \( D_\gamma \) in the usual RT-language are called “suitably defined quantum traces”. The authors

\(^1\) The term “\( \text{dessin d’enfant} \)” or “\( \text{dessin} \)” for short, means exactly the same as “fat graph” or “ribbon graph” and in the following we use all these terms interchangeably.
of \cite{15} provide one of possible consistent definitions of these traces and then derive from this definition some properties. In particular, they show that their definition satisfies recursion relations (MOY-relations), that hypercube formula \cite{4} is invariant w.r.t all Reidemeister moves and that it satisfies skein relation \cite{1}.

In the rest of this section and in the next one we strengthen these statements a little bit, in particular, make them applicable also to virtual knots.

First, observe, that

**Proposition 4.1.** Hypercube formula \cite{7} satisfies skein relation \cite{1} regardless of what $D_\gamma$’s are.

Indeed, consider some black vertex. All summands in \cite{7} split into two big groups: the ones, where the edge, corresponding to this vertex, is kept, and the ones, where it is amputated. Hence, we can write (with self-explanatory notation)

$$H_\bullet = q^{N-1}(-q)H_{\text{kept}} + q^{N-1}H_{\text{amp}}$$ \hspace{1cm} (9)

where $q$-factors, produced by this vertex, are explicitly written down – and thus neither $H_{\text{kept}}$, nor $H_{\text{amp}}$ contain any factors, dependent on the color of our selected vertex.

Now, if this vertex is changed for white one, we can write

$$H_\circ = q^{1-N}(-q^{-1})H_{\text{kept}} + q^{1-N}H_{\text{amp}}$$ \hspace{1cm} (10)

Since all color dependence on the chosen vertex is absent in $H_{\text{kept}}$, $H_{\text{amp}}$, we actually have $H_{\text{kept}} = H_{\text{kept}}$ and $H_{\text{amp}} = H_{\text{amp}}$. This allows us to write

$$A^{-1}H_{\circ} - AH_\circ = (q^{-1} - q) \cdot H_{\text{amp}} = (q^{-1} - q) \cdot H_{||} = z H_{||},$$ \hspace{1cm} (11)

which is exactly the skein relation \cite{1}.

**Example: virtual Hopf link** If we apply the hypercube formula to the virtual Hopf link above, we get

$$H_{v\text{Hopf}} = q^{N-1} \left( D_\bullet - qD_\circ \right)$$ \hspace{1cm} (12)

While $D_\bullet$ poses no questions – it is just a familiar dimension $[N]$ of the unknot, we actually see the new unknown parameter: value of $q$-dimension for the dessin with one vertex and one loop-edge $D_\circ$. Furthermore, we see that this answer satisfies skein relation, in accordance with our general argument – regardless of what we substitute for $D_\circ$.

$$H_{v\text{Hopf}\bullet\rightarrow\circ} = q^{1-N} \left( D_\bullet - q^{-1}D_\circ \right)$$ \hspace{1cm} (13)

$$H_{v\text{Hopf}\bullet\rightarrow||} = D_\bullet$$ \hspace{1cm} (14)

$$A^{-1} \cdot H_{v\text{Hopf}} - A \cdot H_{v\text{Hopf}\bullet\rightarrow\circ} = (q^{-1} - q) \cdot H_{v\text{Hopf}\bullet\rightarrow||}$$ \hspace{1cm} (15)

Thus, we see, that out of two constituents of the skein method – skein relation and Reidemeister moves – the hypercube formula \cite{7} is sensitive only to the first one. It is quite natural to expect that it is the second constituent – Reidemeister invariance – which puts some constraints on the so far completely free parameters $D_\gamma$. As we will see in the following section, this is indeed the case – but in the case of virtual knots these constraints are not sufficient to fully determine all the $D_\gamma$.

We finish this section with a small remark:

**Proposition 4.2.** Hypercube formula is not the most general solution to skein relations

Though formula \cite{7}, with arbitrary $D_\gamma$, solves skein relation \cite{1}, it still has drawbacks. For instance, it automatically respects all the virtual Reidemeister moves \cite{17}. Moreover, it also respects the following additional move – interchange of the usual and virtual crossings, which is not in the list \cite{13}.

\[
\begin{array}{c}
\circ \\
\bullet
\end{array}
\quad = \quad
\begin{array}{c}
\circ \\
\bullet
\end{array}
\]

(16)

This indicates, that it could be possible to write a more relaxed formula, that would be just the generic form of a solution to skein relation. To invent such a formula is a very interesting open question, which, however, is left out of the scope of this paper.
5 Mild assumptions + Reidemeister moves = MOY-relations

So, formula (7) automatically incorporates skein relation. But there is another component of skein method, namely, application of Reidemeister moves to actually unknot the planar diagram after application of skein relations. This Reidemeister invariance would now put some constraints on components of the formula (7), namely, on dimensions $D_\gamma$.

In fact, theorem 5.1 of [13] can be reinterpreted in the following way: suppose dimensions $D_\gamma$ satisfy

- factorization property: $D_\gamma = D_{\gamma_1} \cdot D_{\gamma_2}$ for $\gamma = \gamma_1 \cup \gamma_2$ with $\gamma_1 \cap \gamma_2 = \emptyset$
- normalization: dimension of an isolated vertex is $[N]$, $D_* = [N]$

Then invariance of (7) w.r.t Reidemeister moves implies that $D_\gamma$ are “local” functions of fat-graph $\gamma$ (as Reidemeister moves themselves can be applied to any small part of planar diagram) – they change nicely under certain transformations of the graph, which change only the small piece of it.

These transformations are, in fact, well-known in the literature: they are celebrated MOY-relations [15]. In particular, it is these relations, that inspired Khovanov and Rozansky to invent their categorification of HOMFLY polynomials [19]. However, in [15] these relations are corollaries (Lemmas 2.2, 2.3, 2.4 and 2.5) of an independent (non-recursive) definition of $D_\gamma$. We show ([13]) that for any definition of $D_\gamma$ (satisfying factorization and normalization properties) these relations should be satisfied, to ensure topological invariance.

Armed with hypercube formula (7) and MOY-relations (17)-(21) we are ready to address the question, posed by section 3. First, we rephrase the question in the hypercube language.
Suppose we calculate HOMFLY polynomial for some planar diagram with the help of hypercube method. First we write down the hypercube formula (7) and obtain HOMFLY as a linear combination of dimensions for different fat graphs. Then, using MOY-relations (17)-(21) we gradually simplify these fat graphs, such that, hopefully, in the end they become collections of isolated vertices and we get a Laurent polynomial in $z$ and $A$.

But is it always possible? Thanks to our analysis in section 2 we know, that for dessins, that come from planar diagrams of non-virtual knots there is always at least one way to decompose them, such that result is a number. Moreover, since MOY [15] give explicit formula for these dimensions for planar fat graphs, we know that any two different decomposition paths for a planar fat graph would give the same answer.

But for non-planar dessins we have neither proposition 2.1, nor MOY-formula [15], therefore, whether

- for every non-planar dessin it is possible to obtain a numeric (i.e., not containing other dessins) answer via application of MOY-relations;
- different paths of decomposition of a given dessin would give same answer as a polynomial of simpler dessins (i.e. that there are no additional relations between various $D_i$’s, except MOY)

are open questions – and these are the proper hypercube reformulation of the questions in section 3.

In the following we investigate these questions. Actually we do not have general theoretical arguments, and thus resort to computer checks (which, even if very excessive and persuasive, are by no means sufficient).

Remark 1. Note, that neither the flip-rule ([13], eq.(24)) nor the $1/|N|$-factorization property ([13], eq.(11)) are part of MOY-relations, which are required by topological invariance. They are additional properties, that are most naturally present, if one takes matrix-model point of view, as inspired by [11]). However, an empirical claim of this paper is that they can be safely dropped – this will not break topological invariance of the answer. What does this mean on the matrix-model side of the story is an interesting question for future research.

Remark 2. Also note, that relation (20) was not among the original MOY-relations [15]. This is because the 3rd Reidemeister move with antiparallel orientation of strands (which actually implies (20)) is, topologically, a consequence of parallel 3rd Reidemeister move and antiparallel 2nd Reidemeister move, so MOY decided to omit it. This means that if we did the same, if we had not use (20) to simplify dessins, some dessins would seem indecomposable at first appearance. But then, when considering some dessins with bigger number of edges, we would suddenly see that they have at least two different decomposition paths: one allows to completely simplify this bigger dessin to a number, while the other expresses it though previously indecomposable smaller dessin. Therefore, we would conclude, that this small indecomposable dessin should have some definite numeric value. But, of course, since we are actually using (20) in our search of decompositions, we get the numeric value for the small dessin straight away.

6 Different ways to decompose dessins: tests for topological invariance

The recursion relations (17)-(21) allow to express more complicated dessins (i.e. the ones with bigger number of edges) in terms of simpler ones. However, often there is more than one possible way to do the decomposition.

The simplest example of this phenomenon is a cycle dessin with two vertices and two edges. It can be simplified either with help of $[N - 2]$-rule (19) or $[2]$-rule (18).

\[
\begin{align*}
\begin{array}{c}
\includegraphics[height=0.5cm]{cycle_dessin} \\
\end{array}
\end{align*}
\]

\begin{align}
\begin{array}{c}
\includegraphics[height=0.5cm]{cycle_dessin} \\
= [N - 2] \bullet + \bullet = [N] ([N - 2] + [N])
\end{array}
\end{align}

\begin{align}
\begin{array}{c}
\includegraphics[height=0.5cm]{cycle_dessin} \\
= [2] \includegraphics[height=0.5cm]{cycle_dessin} = [2][N][N - 1]
\end{array}
\end{align}

Thanks to the identity between q-numbers

\[
[N] + [N - 2] = [2][N - 1]
\]

these two ways actually give the same answer. But is it always the case? Do any two different ways to decompose some complicated dessin always give the same answer?

Clearly, for some dessins it is not possible to find a sequence of MOY-moves, that expresses them through dessins with fewer edges, for example, dessin appearing in calculation of HOMFLY for virtual Hopf link
is such a dessin. We call such dessins atomic.

Continuing the analogy, dessins, that can be decomposed, for example

we call molecular.

Now, consider some dessin. All the ways to apply [N-1]-, [N-2]- or [2]-rule to it once should give the same answer (since we assume that value of every dessin is definite). Hence, we require that results of application of there rules are equal. This gives us some relations on dessins with fewer number of edges. We call such relations simple relations.

For example, relation (22) is an example of such relation.

Sometimes, it is also possible to apply [N-3]- or [1]-rule to a dessin. This also gives some relation, but this time between dessins with same number of edges. We call such relations cluster relations.

What we ultimately want to know, is whether these simple and cluster relations give some non-trivial relations on atomic dessins, in particular, whether it follows from them, that

\[
\bigcirc \pm (-) \longrightarrow = -[N][N - 1]
\]

\[
\bigcirc \bigcirc = \frac{1}{[N]} \bigcirc^2 = [N][N - 1]^2
\]

The first equation is a particular case of flip-rule ([13], eq.(24)) and the second is a particular case of \(\frac{1}{[N]}\)-factorization ([13], eq.(11)). We would like to know, whether these, or similar, formulas are corollaries of MOY-relations (17)-(21).

MOY-relations in the hypercube formalism directly correspond to the Reidemeister moves in the skein method. Hence, when we study, which dessins are atomic, which are molecular and what are simple and cluster relations between dessins, we actually learn, what restrictions on possible choice of values of \(D_\gamma\)'s does the topological invariance impose. This is useful, since if an atomic dessin is not restricted by any relations, we can treat it as a new variable in a HOMFLY polynomial. Two candidates for such new variables are dessins \(\bigcirc\) and \(\bigcirc \bigcirc\) (provided we wouldn’t find any of relations (24)).

The extensive search for these (and other) relations, however, is easier said than done. This is because the number of essentially distinct dessins with given number of edges grows very fast (faster, than Catalan numbers). Hence, a computer program is required. We’ve written such a program and briefly outline here the logic of its work.

7 Computer program

- The program works by layers – at each layer it looks at dessins with given number of edges. It is assumed, that expressions for all dessins with fewer edges in terms of atomic dessins are already found.

- It splits dessins on a given layer into clusters – dessins that can be transformed into one another with help of [N-3]-rule and 1-rule (modulo dessins with fewer number of edges). For example, these two dessins on layer 3 are in the same cluster (via [N-3]-rule)

- Then for every dessin it finds all the places where really decomposing rules (i.e. [N-1]-rule, [N-2]-rule or [2]-rule) can be applied to it. It generates equations, that results of all these decompositions should be the same. These results involve expressions for dessins with fewer edges, which, by assumption, we already know. At this point it may happen, that dessin can not be simplified with [N-1]-rule [N-2]-rule or [2]-rule. In this case it is potentially atomic.
Then, it generates equations that say that expressions for all dessins in a cluster are related to each other by 1-rule or \([N-3]\)-rule. For every such equation there may be two cases: at least one dessin involved is potentially atomic or there are no potentially atomic dessins. In the latter case we have some equation on dessins from lower layers. In the former case we express one of potentially atomic dessins using this equation and continue.

There are couple of technical details that are important for realization of the algorithm:

- At every level we need a persistent naming (numbering) of dessins, because we want to have a way to save expressions for them in terms of atomic dessins and then fetch these expressions.
- Since some relations may at first look non-trivial, but then completely simplify to zero due to identities for \(q\)-numbers, the CAS is needed to do this polynomial algebra. We use Wolfram Mathematica. Partly because it’s syntax and idioms are so close to Common Lisp, in which the rest of our program is written.

This program is part of our CL-VKNOTS program and can be found on Knotbook [16]; as is customary for all the source code nowadays, it can also be found on GitHub [20].

8 The results

The more extensive statistics of the program’s work is in the appendix [15]. Here we present the most striking and important features of it.

Up to and including 7 crossings we are not finding any nontrivial relations between atomic dessins. Number of atomic dessins with 0, 1, 2, 3, 4, 5, 6 and 7 edges is equal, respectively, to 0, 1, 1, 5, 15, 91, 784 and 9160. Here are first few of them, for small number of edges.

| # edges | dessins |
|---------|---------|
| 1       | ✐      |
| 2       | ✐ ✐    |
| 3       | ✐ ✐ ✐ ✐ ✐ ✐ ✐ |
| ...     | ...    |

All these dessins supposedly can be used as new independent variables in HOMFLY polynomial. Furthermore, thanks to MOY-formula [15] we know that for planar dessins (i.e., non-virtual knots) these additional variables never occur. Therefore, if HOMFLY polynomial for some planar diagram (involving virtual crossings) contains these parameters, then this diagram corresponds to essentially virtual knot – virtual crossings can not be eliminated from this diagram by Reidemeister moves. Unfortunately, if HOMFLY does not contain these additional parameters, it does not mean the knot is non-virtual (as appendix A shows).

9 Imposing factorization of composites

As an example in Appendix A illustrates, inclusion of all our new parameters breaks factorization property of fundamental HOMFLY for composite knots. This property does not follow from topological invariance, still it is common for all knot polynomials, studied so far – so one can wish to preserve it.

For this purpose it is sufficient to impose additional relation on dessins \(-1/[N]\)-decomposition rule of [13].

\[
[N] \cdot D \left( \begin{array}{c}
\gamma_1 \\
\gamma_2
\end{array} \right) = D \left( \begin{array}{c}
\gamma_1 \\
\bullet
\end{array} \right) \cdot D \left( \begin{array}{c}
\bullet \\
\gamma_2
\end{array} \right)
\]

(25)

In other words, if a dessin has a vertex, which we can cut so that it becomes disconnected, then we can express dimension of the whole through dimensions of parts.

This considerably reduces the number of dessins, that are independent – to
but does not eliminate them completely. Moreover this number still grows with the number of edges, so the number of new parameters still seems to be infinite.

10 Conclusion

In this paper we continued our study of hypercube formalism, applied to virtual knots and links. We tried to drop as many assumptions on the form of the answer as possible, keeping only those, that are dictated by skein relation and topological invariance. In particular, we dropped flip-rule and $1/[N]$-decomposition rule of [13].

This way, we obtained surprising empirical results. Namely, dimensions of many hypercube vertices are not fixed by topological invariance and can, therefore, be treated as new variables in HOMFLY polynomial. Inclusion of these parameters breaks some properties, which are well-known for HOMFLY polynomials for non-virtual knots but are subject of discussion for virtual ones, such as factorization of HOMFLY polynomial for composite knots. Moreover, even imposing factorization of composites as additional requirement (despite it does not follow from topological invariance), we only diminish the number of new parameters, while many still seem to survive.

These parameters never appear in HOMFLY polynomials for non-virtual knots and therefore provide sufficient tool to distinguish essentially virtual knots. Also, no properties of HOMFLY polynomials for non-virtual knots are broken as polynomials remain the same.

For $N = 2$, i.e. for Jones polynomials, a more powerful approach to knot polynomials is possible. It uses different skein relation [6] that, rather than changing the crossing for inverse one, completely replaces crossing by its two different Morse resolutions. When applied to virtual knots (what is actually done in [6]) this approach has no ambiguities, since after all non-virtual crossings have been eliminated using this Jones skein relation, all diagrams always completely unknot. This contrast between HOMFLY-skein and Jones-skein relations is not surprising: they are already very different for non-virtual knots. While the latter one allows for recursive calculation of its polynomial, the former one needs an additional search for unknotting crossings, which makes calculation an enumerative method, not a recursion.

Since this paper mixes presentation of general theory with empirical results, i.e. facts with assumptions, the logic is not fully consistent – what is, however, natural for the subject which is not fully understood. As we do not have general proofs, we resort to conjectures and checks. Checks are necessarily done with the help of computer, as the problem quickly explodes combinatorially. We check our conjectures for planar diagrams with up to and including 7 non-virtual crossings and arbitrary number of virtual ones.

We formulate two conjectures, minimal and maximal, which can be briefly stated as follows:

**Conjecture 10.1.** At least one new parameter, consistent with topological invariance, can be put into HOMFLY polynomial for virtual knots. It is the value of $q$-dimension of the dessin with one vertex and one edge: $\bullet$. This parameter can not be fixed, even if factorization of composites is requested.

**Conjecture 10.2.** An infinite number of free parameters, which are values of $q$-dimensions for dessins, that can not be decomposed using $[N-1]$-, $[N-2]$-, $[2]$-, $1$- and $[N-3]$-rules can be put into HOMFLY polynomial for virtual knots, again preserving its topological invariance. Many of them are consistent with (i.e. not fixed by) factorization of composites.

The difference is that empirical evidence in support of the former conjecture looks overwhelming, while the checks of the latter one are a little less exhaustive.

Extension of these conjectures to colored HOMFLY of virtual knots and clarifying of their relation to Chern-Simons/RT approach to HOMFLY polynomials are important subjects of future research.
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A  Kishino knot and its relatives

Kishino knot (4.55 & 4.56 in the table [21]) is a composite of two virtual unknots, which, however, is not topologically trivial.

\[
\begin{array}{c}
\text{A Kishino knot and its relatives} \\
\end{array}
\]

First we calculate HOMFLY polynomials for it, and related knots, using well-developed methods of [11]. Then, we use method of this paper to see, how it changes the situation.

A.1 Gradual quantization (matrix-model) method

Primary and full classical hypercubes for the main diagram

\[
\begin{array}{c}
\text{calculated by the rules of [12], are} \\
1 \\
2 3 2 \\
2 3 2 \\
1 \\
2 3 2 \\
2 3 2 \\
\end{array}
\]

\[
\begin{array}{cccc}
1 & 2 & 3 & 2 \\
2 & 3 & 2 & 2 \\
2 & 3 & 2 & 2 \\
1 & 1 & N & N \\
2 & 3 & 2 & 2 \\
2 & 3 & 2 & 2 \\
2 & 3 & 2 & 2 \\
1 & 1 & N & N \\
\end{array}
\]

\[
\begin{array}{c}
\text{Quantization in this case is straightforward and the full hypercube is:} \\
\end{array}
\]

\[
\begin{array}{cccc}
-[N] & [N] & [N] & [N] \\
-[N] & [N] & [N] & [N] \\
-[N] & [N] & [N] & [N] \\
-[N] & [N] & [N] & [N] \\
\end{array}
\]

\[
\begin{array}{cccc}
-2[N][N] & [N][N] & [N][N] & [N][N] \\
-2[N][N] & [N][N] & [N][N] & [N][N] \\
-2[N][N] & [N][N] & [N][N] & [N][N] \\
-2[N][N] & [N][N] & [N][N] & [N][N] \\
\end{array}
\]

Starting from the main (Seifert) vertex, marked by a single box, we obtain for

\[
\begin{array}{c}
\text{(26)} \\
\end{array}
\]
\[ q^{4N-4}[N] \left\{ 1 + 4q[N-1] + 4q^2[N-1]^2 - 2q^3[N-1] - 4q^3[2][N-1] + q^4[2]^2[N-1]^2 \right\} = \]
\[ = q^{4(N-1)}[N] \left( 1 - q^2(q - q^{-1})[N-1] \right)^2 = [N] \left( q^{N+1} + q^{2N-2} - q^{3N-1} \right)^2 = [N] \cdot \left( v_{[2,1]}^2 \right)^2 \]  
(27)

as one should expect for a composite of the two virtual knots 2.1.

Similarly for

\[ q^{-2}[N] \left\{ -2[N-1] - q \left( 2[2][N-1]^2 - 2[N-1] \right) + q^2 \left( 1 + 2^2[N-1]^2 + 4[N-1]^2 \right) - \right. \]
\[ -q^3 \left( 2[2][N-1]^2 - 2[N-1] \right) - q^4[2][N-1] \right\} = [N] \left( 1 - 2[N-1](q - q^{-1}) \right)^2
\]
\[ + [N-1]^2(q - q^{-1}) \right) = [N] \left( 1 - q^2(q - q^{-1})[N-1] \right)^2 = [N] \cdot v_{[2,1]}^2(q) \cdot v_{[2,1]}^2(q^{-1}) \]  
(28)

– what is again the expected answer for a composite of 2.1 and its copy with inverted orientation.

Thus it comes with no surprise that the Kishino knot itself, which is a composite of two virtual unknots, has a trivial HOMFLY. Indeed, taking the triple-boxed vertex as initial, we get for either

\[ = \]
(29)

or

\[ = \]
(30)

\[ q^{-2}[N] \left\{ [N-1]^2 - q \left( -2[N-1] + 2[2][N-1] \right) \right\} + q^2 \left( 1 + 2[2][-][N-1] + 2[N-1]^2 \right) - q^3 \left( 2(-)[-][N-1] + 2[2][N-1] \right) + q^4[2][N-1]^2 \}
\[ = [N] \]  
(31)

A.2 Axiomatic (MOY-relations) method

Let’s now calculate HOMFLY for Kishino knot and related knots using techniques discussed in this paper. We start with the knot in the Seifert vertex (picture (26)). Corresponding dessin d’enfant is

\[ \]  

hence HOMFLY polynomial, according to the amputation formula (7), is

\[ \]
\[ H_{\text{Seifert}} = q^{4N-4} \left( q^4 - 4q^3 + 4q^2 + 2q - 4q \right) + \left( [N] + q^2 \left( q - \frac{1}{q} \right) \right) \]

\[ = q^{4N-4} \left( [N] + 2q^2 \left( q - \frac{1}{q} \right) \right) \]

On the other hand, the planar diagram of a virtual trefoil (2.1 in the table [21]), has dessin d’enfant

\[ = \left( [N] + 2q \left( q - \frac{1}{q} \right) \right) \]

hence for its HOMFLY polynomial we get

\[ H_{2.1} = q^{2N-2} \left( [N] + q^2 \left( q - \frac{1}{q} \right) \right) \]

so the factorization property for composites \( H_{\text{Seifert}} = 1/[N]H_{2.1}^2 \) is no longer true, unless, of course, one imposes by hand (we have not found any evidence, that it’s required by the topological invariance) that

\[ \bullet \quad \bullet = \frac{1}{[N]} \bullet^2 \]

The two versions of Kishino knot (pictures (29) and (30)) have the following dessins d’enfant (dotted edges correspond to white vertices of planar diagram)

The combinatorics of \( q \)-factors in amputation formula (7) is the same for both of them and reads (here we omit difference between solid/dotted lines, since we’ve already written out all \( q \)-weights explicitly)

\[ H_{\text{Kishino}} = \left( \right) \]

\[ = \left( [2]^2 \bullet \bullet - 2[2]^2 \bullet \bullet + [2]^2 \bullet \bullet + 2[2] \bullet \bullet - 2[2] \bullet \bullet \right) \]

Thus our new independent parameters cancel from HOMFLY polynomial. This shows that appearance of these new parameters is sufficient for a knot to be essentially virtual, but not necessary.

### B Results of the program’s work

Even with computer program we are able to get to dessins with only up to and including 7 edges. This is because the number of distinct (connected) dessins with given number of edges grows as

| # edges | 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  |
|---------|----|----|----|----|----|----|----|----|
| # dessins | 1  | 2  | 5  | 20 | 107| 870| 9436| 122840 |
The number of relations, that come from application of \([N-1]\)- [N-2]- or [2]-rules to different places of the same \textit{dessin} is, respectively

| # edges | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|---------|---|---|---|---|---|---|---|---|
| # simple rels | 0 | 1 | 4 | 14 | 86 | 729 | 8385 | 113399 |

The number of \textit{dessins}, decompositions of which give the given number of equations

1: 1 \rightarrow 1
2: 1 \rightarrow 5, 2 \rightarrow 2
3: 1 \rightarrow 32, 2 \rightarrow 16, 3 \rightarrow 6, 4 \rightarrow 1
4: 1 \rightarrow 265, 2 \rightarrow 136, 3 \rightarrow 49, 4 \rightarrow 10, 5 \rightarrow 1
5: 1 \rightarrow 2793, 2 \rightarrow 1588, 3 \rightarrow 580, 4 \rightarrow 140, 5 \rightarrow 22, 6 \rightarrow 1
6: 1 \rightarrow 36068, 2 \rightarrow 20996, 3 \rightarrow 8175, 4 \rightarrow 2162, 5 \rightarrow 385, 6 \rightarrow 39, 7 \rightarrow 1

However, all these relations actually \textit{simplify to zero}, thanks to identities between \(q\)-numbers (i.e. coefficient in front of every monomial in atomic \textit{dessins} is separately equal to zero).

The number of cluster relations (i.e. the ones which relate \textit{dessins} by \([N-3]\)- or 1-rule) is

| # edges | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|---------|---|---|---|---|---|---|---|
| # cluster rels | 0 | 0 | 2 | 18 | 250 | 3415 | 49316 |

Among them, the following number contains potentially atomic \textit{dessins} and hence is used to express these potentially atomic \textit{dessins} through sub-\textit{dessins}

| # edges | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|---------|---|---|---|---|---|---|---|
| # rels | 0 | 0 | 0 | 3 | 46 | 630 | 8631 |

All other cluster relations connect two \textit{dessins}, each of which is independently simplifiable by \([N-1]\)- [N-2] or [2]-rule, and hence in principle could’ve given some relation between sub-\textit{dessins}, but all of them \textit{simplify to zero}.

The number of \textit{dessins}, that are expressed either through sub-\textit{dessins}, or, via cluster relations, through \textit{dessins} on the same level, is

| # edges | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|---------|---|---|---|---|---|---|---|---|
| # molecular dessins | 1 | 1 | 4 | 15 | 92 | 779 | 8652 | 113680 |

Hence, the number of atomic \textit{dessins} at each level (i.e. number of potentially atomic minus the ones which were expressed using cluster relations on this level) is

| # edges | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
|---------|---|---|---|---|---|---|---|---|
| # atomic dessins | 0 | 1 | 1 | 5 | 15 | 91 | 784 | 9160 |

In result, it looks like we have a lot of atomic, indecomposable \textit{dessins} and their number grows fast with the number of edges. However, it still may be the case that we are not looking into the structure deep enough. Since number of different \textit{dessins} with a given number of edges also grows fast, it may be that at certain point (say, 20 edges) they start to provide many constraints on values of \textit{dessins} with lower number of edges. Say, 20-edge \textit{dessins} could constrain completely 2-edge \textit{dessins}, 21-edge \textit{dessins} – 3-edge \textit{dessins} and so on. Since there is no hope to address the issue by brute force (due to combinatorial explosion), a more intelligent method is needed. This is, however, beyond the scope of the present, largely empirical, investigation.
In this section we develop hypercube construction slightly differently. The difference is that instead of projector onto antisymmetric representation $P_{[1,1]}$ central role is played by projector onto symmetric representation $P_{[2]}$. If one does not keep in mind, that HOMFLY admit Khovanov-Rozansky categorification, this view should be no worse than traditional, antisymmetric one. This allows us to show, that some choices, that seem canonical and rigid in the antisymmetric picture are not so fixed.

Namely, we can start with formula for $R$-matrix through symmetric projector, instead of symmetric one

\[ R = |2|P_{[2]} - qI \otimes I \]

\[ R^{-1} = |2|P_{[2]} - \frac{1}{q}I \otimes I \]

Then, instead of hypercube formula (7) one gets the following analogous formula

\[ H_L^{\square} = |q^N(-q)^{n_\bullet(\Gamma_{Lc}) - n_\circ(\Gamma_{Lc})}\sum_{\gamma \subseteq \Gamma_{Lc}} \left(-\frac{1}{q}\right)^{n_\bullet(\gamma)-n_\circ(\gamma)} \cdot D_\gamma(q, N) \]

where now edges of fat graphs correspond to $|2|P_{[2]}$.

Invariance with respect to Reidemeister moves, together with factorization and normalization properties (see section 5) lead to the following relations on these, symmetric, fat graphs

\[ \begin{align*}
\text{circ} & \quad \simeq \quad [N + 1] \\
\text{triangle} & \quad \simeq \quad |2| \\
\text{double triangle} & \quad \simeq \quad [N + 2] \\
\text{square} & \quad \simeq \quad [N + 3] \\
\text{double square} & \quad \simeq \quad [N + 3]
\end{align*} \]

Again, using only these relations we are not able to uniquely determine numeric value of every dimension $D_\gamma$. In particular, we do not know, what is the value of $\bullet$. But what is the natural value of it?

\[ \text{double square} \quad \simeq \quad [N + 3] \]
At \( q = 1 \) we can use matrix model and see, that the value is \( N(N+1) \). So, the most natural guess would be that the quantum value is \( [N][N+1] \).

But if then, based on this value, we would try to calculate the value of the old loop dessin in antisymmetric picture (using \([1] \otimes [1] = [1, 1] + [2]\)) we would get

\[
\bigcirc_{\text{antisym}} = [2][N] - \bigcirc_{\text{symm}} = [2][N] - [N][N+1] \quad (43)
\]

Classically, at \( q = 1 \) this value is still equal \(-N(N-1)\), but the quantization is different from \(-[N][N-1]\), which seems most natural from “antisymmetric” point of view.

Vice versa, if we take the value \(-[N][N-1]\) for antisymmetric loop diagram and calculate symmetric one based on it, we would get \([2][N] + [N][N-1]\) – different quantization, than the one that seems natural from “symmetric” point of view.

Since symmetric and antisymmetric points of view are so interchangeable (related just by transposition of the Young diagram) there is a priori no reason to prefer one over the other. But choosing one of them would mean that flip-rule (24) is not true, whereas it holds for the other. This is an indication that the flip rule can not be a consequence of topological invariance. It is some extra structure, that is related to preferring one channel of representation theory over another, in other words, to symmetry breaking.

This is where we would like to stop this very intriguing inquiry into representation theory, for now.

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