Preface to the Special Volume on the SAT 2005
Competitions and Evaluations

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1. From a single event to a special volume: History of the SAT competitions

In 2005 there were numerous “competitive events” in the area of automated reasoning (in the broad sense): The Tenth CASC competition[23], the First Satisfiability Modulo Theory Competition[3], the First CSP competition, the Fourth SAT competition, the Third QBF evaluation and the First Pseudo Boolean evaluation, etc.

Starting such an event is suitable for promoting a common input format and to build a repository of benchmarks in that format: this was the motivation behind the SMT and CSP competitions and the PB evaluation. The First QBF evaluation was also organized two years ago in that spirit and the renewal since then allowed to impose the QDimacs input format and to increase both the number of solvers and benchmarks available each year.

More mature events, such as the CASC competitions, are a bit different: both the pool of benchmarks (TPTP) and the pool of solvers are quite stable and the competitions allow to track the progress of new versions of those solvers on a well studied set of benchmarks.

The SAT competitions are again different: the first SAT competition is older than CASC (1992, [6]), a common input format exists since the Second Dimacs challenge (13 years ago [12]), many SAT benchmarks are available (from a repository such as SATLIB [11] or from academic or private companies web sites) in that format. Because recent SAT benchmarks are really huge (several hundreds of MB each) and others cannot be freely distributed, it is difficult to think about a common repository similar to the TPTP where all the SAT benchmarks could be available. Furthermore, the number of available SAT solvers is growing each year and makes it difficult to maintain up-to-date initiatives such as SATEX [22]. As a consequence, the SAT competition should be seen as a yearly snapshot of a subset of current SAT solvers on a subset of the available SAT benchmarks.

The four consecutive SAT contests have brought the community a large amount of new solvers and benchmarks, and a place where both solver designers and users can meet. The main difference between the initial SAT 2002 competition and the current SAT 2005
Table 1. SAT Competition Overall Progress. The \textit{“†”} mark means that three launches were attempted for each randomized solvers, with the same timeout.

|          | 2002 | 2003 | 2004 | 2005 |
|----------|------|------|------|------|
| #Solvers | 28   | 34   | 55   | 43   |
| #Submit. | 19   | 22   | 27   | 22   |
| Timeout (s) 1st | 2,400† | 900  | 600  | 1,200 |
| Timeout (s) 2nd | 21,600† | 7,200 | 2,400 | 6,000/12,000 |
| Total CPU (days) | 750  | 522  | 516  | 1,230 |
| Judges    | –    | J. Franco, H. van Maaren, T. Walsh | F. Bacchus, J. Marques-Silva, H. Kleine Buning, A. van Gelder, A. Biere, O. Kullmann |

competition is that SAT solvers are currently being used in various areas, in academia or in the industry, because they can handle some benchmarks with millions of variables and clauses. In the SAT 2005 competition, this is reflected by the variety of benchmarks in the industrial category, and the participation of companies such as IBM and Intel.

1.1 What’s new in the SAT 2005 competition

Over the years, new rules were introduced in order to ensure that the community may benefit from the event. We mainly focused on the “anti-black-box” rule in the 2004 contest, but we relaxed a bit that rule this year. It was decided to mimic CASC and to have two divisions in the SAT 2005 competition: the competition division and the demonstration division. The competition division is the place where solvers, available in source for the research community, can compete and from which winners are selected. The demonstration division is the place where solvers available in binary form, or other solvers whose presence could benefit the research community, are allowed. There was no winner in that second category. At most 3 variants of the same solver (same submitter) were allowed in the competition division. Additional submitted solvers were moved to the demonstration division (the submitters chose which solver ran in which division). This year the three judges were Armin Biere, Allen van Gelder and Oliver Kullmann.

For the first time, we had less solvers submissions than the previous year. We had 43 solvers (see table 1 for details over the years), and we selected 1657 benchmarks over the three categories, 390 random, 675 crafted and 592 industrial (296 original + 296 shuffled). The total amount of CPU time was 1230 CPU-Days on two different clusters of 18 computers each, 680 days of CPU-time for the first stage and 550 days of CPU-time for the second one.

Practically speaking, the hardware was composed of 16 Athlon 1800+ with 1GB RAM, provided by the LRI, Université de Paris-Sud, and with 8 Athlon 1800+ with 2 GB RAM, provided by the LINC Lab, Department of ECECS, University of Cincinnati. Computers were running GNU Linux (RH flavor); Solvers were compiled with GCC 3.3.5.; and the Java solver was launched on Sun Java 1.5.0.02 JVM.
2. The new scoring scheme

The importance of the international SAT competition has grown to being an awaited event in the community. The major impact of being ranked among the best solvers is beneficial both for academic and industrial competitors. As it was stated in our previous reports, our ranking was until now based on a series-basis. One main problem of our scoring scheme (see previous competition report for details [21, 4, 5]) was that all solvers in the second stage were awarded the same score (they solved the same number of series). The final ranking of the second stage was often based on the simple number of solved benchmarks, and the scores were very tight (on some categories, the winner solved only one additional benchmarks than the second one). As an additional argument, because solvers may be used with very large CPU-timeout, it is now important to consider the CPU-time of solvers. A fast solver must be preferred to another one if they both solve the same amount of benchmarks, because the faster one may be efficiently embedded in a real-world application where it will have to scale-up well.

2.1 Design Objectives

One key idea behind the SAT competition is to award a solver that is good on a wide range of SAT instances. In the previous year of the competition, this was implemented using a scoring scheme that ranked the solvers with a tiered system: First, the solvers were ranked by being able to solve some instance in a highest number of different series. Ties were then broken using the total number of benchmarks solved. Unfortunately, in this system there is no difference between solving a benchmark solved by all solvers or one solved by only a few solvers. The same applies to series too.

Another key idea of the competition was to focus on solvers that are the only ones to solve some benchmarks: in the SAT and CASC competitions, those solvers are called state-of-the-art contributors (abbreviated SOTAC). In the previous scoring scheme, the solvers did not benefit directly for being SOTAC in their category, even though SOTAC solvers were usually among the top ranked solvers.

Third, the time needed to solve a given benchmark also needs to be considered. While the CPU time was indirectly used for scoring the solvers in the previous years of the SAT competitions, by using a fixed timeout per benchmark, there was no way to discriminate among the solvers able to solve a given benchmark within that timeout.

Furthermore, the second stage ranking was based only on the number of benchmarks solved during the second stage, among those benchmarks that had not been solved by any solver during the first stage. This criterion is based on very strong assumptions:

- The remaining benchmarks are representative of the initial set of benchmarks.
- The solvers will behave in the second stage in a way similar to the first stage.

However, these assumptions did not necessarily hold. Although it is likely that the winners of the previous competitions could have been declared winners using various scoring schemes, nevertheless, the rankings of the remaining top solvers could have changed a lot.

The scoring scheme used for the SAT 2005 competition is designed to address these issues. It incorporates these features:
It gives more credit for solving hard benchmarks than solving easy ones.

- It gives more credit for solving a benchmark fast.

- It gives extra credit for each series solved.

- It stabilizes the rankings of the solvers at the end of the competition.

While the scoring scheme was designed on a purely theoretical basis, the results of the SAT 2005 Competition indicate that the new scoring scheme meets its expectations in practice.

2.2 The Purse-Based Scoring System

The implemented scoring plan works as follows. A run is defined to be the execution of one solver on one benchmark instance, or problem. Each run is allocated a certain amount of CPU time. If the solver succeeds, timeUsed records the time.

For SAT 2005, there are three categories of benchmark, INDUSTRIAL, CRAFTED, and RANDOM. Within each category, there are several specialties, such as SAT, SAT+UNSAT, UNSAT, and CERTIFIED-UNSAT. The scoring system is applied separately within each combination of category and specialty.

Each problem has a solution purse, which is divided equally among all competition solvers that solve the problem. For SAT 2005, all problems have the standard solution purse (stdP).

Each problem has a speed purse, which is divided unequally among all competition solvers that solve the problem. The speed purse is a fixed multiple (spdM) of the solution purse for all problems in the entire competition; it gives a weighting between solving and speed.

The formula to divide the speed purse of a problem is the following, where \( p \) is problem-id and \( s \) and \( i \) are solver-ids, times are in seconds, and 10,000 is an arbitrary scale factor.

\[
speedFactor(p, i) = \begin{cases} 
\frac{10000}{1 + timeUsed(p, i)} & \text{if } i \text{ solved } p; \\
0 & \text{if } i \text{ did not solve } p. 
\end{cases} 
\]  
(1)

\[
speedAward(p, s) = \frac{speedPurse(p) \times speedFactor(p, s)}{\sum_i speedFactor(p, i)} 
\]  
(2)

Thus, the speedAward is pro rata by speedFactor.

The series purses reward breadth of application. Each series (within specialty within category) has a series purse, which is divided equally among all competition solvers that solve at least one problem in the series. If no solver solves any problem in a certain series, its series purse is not distributed.

For SAT 2005, all series containing 5 or more benchmark instances have the same series purse, which is a fixed multiple (serM) of the standard solution purse. (Recall that scoring is separately applied within each combination of category and specialty, e.g., SAT within RANDOM, or SAT+UNSAT within CRAFTED.) All series containing 4 or fewer benchmarks have the same series purse, which is a fixed multiple (serM / 3) of the standard solution purse.
The coefficients and multiples for SAT 2005 are:

\[ stdP = 1000.0; \quad spdM = 1.0; \quad serM = 3.0. \]

3. The SAT 2005 settings

3.1 Submitted Solvers

A 2-pages description of solvers is available on the competition web site. Briefly, the submitters to the competition division were: Anbulagan [2] (dew-satz-1a, dew-satz-1b, dew-satz-1c); Domagoj Babic (hsat.1, hsat.5, hsatrr); Armin Biere (compsat), pages 201–208; Gilles Dequen, Olivier Dubois [7] (kcnfs-2004); Niklas Een, Niklas Sörensson [8] (SatELiteGTI); Zhaohui Fu [15] (zchaff); Roman Gershman [9] (haifasat, haifasat2); Marijn Heule, Mark Dufour, Joris van Zwieten, Hans van Maaren (pages 47-59, march-d1); Holger Hoos, Dave Tompkins [24] (adaptnovelty); Daniel Le Berre, Mederic Baron, Geoffrey Bourgeois (sat4j.jar); Chu Min Li, Wen Qi Huang [13] (g2wsat); Feng Lu, Kai Yang, Kwang-Ting Cheng (csat); Yogesh Mahajan, Sharad Malik, Lintao Zhang, Zhaohui Fu [15] (zchaff-rand); Alexander Nadel (jerusat1.31-a, jerusat1.31-b); Richard Ostrouski [18] (lsatv1.1, wilsatv1); Duc Nghia Pham, Anbulagan (ranov, rpaws10, rrsaps); Steven Prestwich [19] (vw); Niklas Sörensson, Niklas Een [8] (minisat-static); Ivor Spence (tts-3-o); Dave Tompkins, Holger Hoos, Frank Hutter [24] (saps); Daniel Vallstrom, (vallst.sh)

In the demonstration division, we had: Anbulagan (dew-satz-1d, dew-satz-1e); Gilles Dequen, Olivier Dubois (kcnfs); Niklas Een, Niklas Sörensson (satelite-release); HoonSang Jin, Fabio Somensi (circusa, circusb, circusd) demonstration; Raihan Kibria, Niklas Een, Niklas Sörensson (midisat-static); Alexander Nadel, Ziyad Hanna (eureka-a, eureka-b, eureka-c); Duc Nghia Pham, Anbulagan (rpaws40, rpaws5)

No portfolios (such as satzilla in the previous contest) were submitted this year and kcnfs-2004 (same binary as last year winner) was resubmitted by its authors (Gilles Dequen, Olivier Dubois).

3.2 Benchmarks

The random category only contains pure, uniform, non-biased random k-SAT instances, not forced to be SAT or UNSAT. Oliver Kullmann provides full details and analysis of this category pages 61-102.

Crafted benchmarks contained all benchmarks that were neither random nor industrial. Previous tricky benchmarks that were simple but known as needing an exponential number of resolutions (like Urquhart or XOR-chains problems) were mostly taken out last year by lsat and not submitted again this year. Benchmarks that looked like random (for instance forced SAT random benchmarks) were also pushed into this category. As a matter of fact, crafted benchmarks represented the largest set of benchmarks this year (675 benchmarks). We had benchmarks of very different kinds and we shuffled them all of them for the contest. Armin Biere submitted LinvRinv benchmarks; Matti Järvisalo submitted benchmarks based on 3-Regular Graphs (see details pages 27–46); Ashish Sabharwal submitted counting, ordering and pebbling problems; Inês Lynce submitted social golfer problems [14]; Volker Sorge submitted algebraic benchmarks problems [16]; Klas Markström submitted a
family of Eulerian graphs (pages 221-228) and Oliver Roussel submitted PHNF encoding [20] of last year contest, medium hardness, benchmarks.

In the industrial category, new formal verification benchmarks from IBM were submitted by Emmanuel Zarpas. He discusses the choice of the benches pages 229-236. Miroslav Velev submitted previously-known benchmarks (VLIW-SAT (2.0 and 4.0), VLIW-UNSAT 2.0 and Liveness UNSAT 2.0) [25]. Sanjai Narain submitted VPN models from Alloy [17]; François Grieu submitted VPMC inversion problems and Frédéric Maris submitted Planning problems.

3.3 First stage results

The results of the first stage are detailed in Table 2 for the random category, Table 3 for the crafted category and Table 4 for the industrial category.

3.4 Second stage results: The winners

For the second stage, the judges were asked to select the Top-N solvers in each category (in each category, N may be different). We then used a significantly larger cpu-time (6,000 seconds for the RANDOM and CRAFTED categories; and 12,000 seconds for the INDUSTRIAL category) to relaunch the selected solvers on all the benchmarks on which they failed during the first stage. The Table 5 gives the results for the second stage of the competition for random benchmarks. The Table 6 gives the results for the second stage of the competition for crafted benchmarks, and the Table 7 gives the results for the second stage of the competition for industrial benchmarks. More information, including various figures and the details of all runs can be found on the competition web site: http://www.satcompetition.org/2005/.

4. Introducing this special volume

The special volume contains two types of contributions: articles and research notes. The articles are expected to reach a wider audience than research notes. All the contributions have been reviewed in the same way. Considering the specific topic of the special volume, the SAT competition and the Quantified Boolean Formulae (QBF) and Pseudo Boolean (PB) evaluations organized as satellite events of the SAT conference, the contributions are either reports of events (pages 61–164, 229), description of technologies behind solvers (pages 1, 47, 65, 191–219), or analysis of new classes of benchmarks (pages 27, 221).

One of the big winner of the SAT 2005 competition and evaluations is certainly MiniSat. It is the second solver in numerous categories in the SAT competition, and the first when used with the SatELite preprocessor. In conjunction with a pseudo boolean preprocessor, it also shown a great potential in the first pseudo boolean evaluation. Niklas Eén and Niklas Sörensson are presenting in the first paper, Translating Pseudo-Boolean Constraints into SAT, pages 1–26, the original approach taken by MiniSat+: translating the pseudo boolean constraints into clauses, on which the powerful minisat SAT solver can be launched.

Current SAT solvers are now able to handle larger and larger benchmarks, scaling up to millions of variables and clauses. In the second paper, Hard Satisfiable Clause Sets for Benchmarking Equivalence Reasoning Techniques, pages 27–46, Haari Haanpää et al. are interested in finding a model to generate small, hard, satisfiable benchmarks. They
introduce a new family of benchmarks, regular XORSAT, that was submitted to the SAT 2005 competition, crafted category.

If conflict driven clause learning solvers are suitable for solving industrial benchmarks, lookahead solvers with sophisticated heuristics and deduction processes are another appealing approach for solving both random $k$-SAT and crafted benchmarks. Marijn Heule and Hans van Maaren introduce in the third paper, *March_dl: Adding Adaptive Heuristics and a New*
Table 3. First stage results, CRAFTED category

(A) SAT+UNSAT

| Solver        | Score  | #Solved | Sat | Unsat |
|---------------|--------|---------|-----|-------|
| satelitegti   | 79,069 | 173     | 145 |       |
| vallst-sh     | 60,869 | 195     | 130 |       |
| tts-3-0       | 56,951 | 20      | 78  |       |
| march-dl      | 55,428 | 182     | 111 |       |
| zchaff-rand   | 52,064 | 195     | 119 |       |
| minisat-static| 51,536 | 163     | 136 |       |
| csat          | 49,476 | 158     | 113 |       |
| jerusat1-31-a | 43,555 | 167     | 109 |       |
| hsat-1        | 41,719 | 158     | 105 |       |
| hsat-5        | 41,213 | 157     | 105 |       |
| zchaff        | 37,601 | 169     | 97  |       |
| haifasat      | 33,840 | 156     | 95  |       |
| jerusat1-31-b | 32,660 | 167     | 100 |       |
| hsatrr        | 31,482 | 111     | 66  |       |
| haifasat2     | 29,790 | 153     | 96  |       |
| compsat       | 23,888 | 143     | 83  |       |
| dew-satz-1b   | 21,250 | 125     | 67  |       |
| sat4j.jar     | 21,719 | 136     | 73  |       |
| dew-satz-1a   | 21,063 | 127     | 66  |       |
| dew-satz-1c   | 19,868 | 124     | 63  |       |
| lsatv1-1      | 19,055 | 61      | 37  |       |
| vw            | 17,047 | 130     | 0   |       |
| wllsatv1      | 11,912 | 86      | 50  |       |
| rrsaps        | 11,312 | 146     | 0   |       |
| g2wsat        | 10,393 | 128     | 0   |       |
| ranov         | 9,561  | 134     | 0   |       |
| kcnfs-2004    | 8,269  | 49      | 34  |       |
| rpaws10       | 7,875  | 101     | 0   |       |
| saps          | 5,221  | 97      | 0   |       |
| adapttnovelty | 5,130  | 87      | 0   |       |

(B) SAT

| Solver        | Score  | #Solved | Sat | Unsat |
|---------------|--------|---------|-----|-------|
| satelitegti   | 38,063 | 173     | 145 |       |
| march-dl      | 31,656 | 182     | 105 |       |
| vallst-sh     | 29,366 | 195     | 105 |       |
| minisat-static| 23,913 | 163     | 158 |       |
| hsat-1        | 23,502 | 158     | 105 |       |
| hsat-5        | 22,894 | 157     | 105 |       |
| zchaff        | 21,004 | 169     |     |       |
| jerusat1-31-a | 20,919 | 167     | 109 |       |
| haifasat      | 20,078 | 156     |     |       |
| csat          | 19,257 | 158     |     |       |
| zchaff-rand   | 19,875 | 159     |     |       |
| jerusat1-31-b | 18,887 | 167     |     |       |
| hsatrr        | 18,487 | 111     |     |       |
| vw            | 17,869 | 130     |     |       |
| haifasat2     | 16,959 | 153     |     |       |
| compsat       | 14,952 | 143     |     |       |
| sat4j.jar     | 12,395 | 136     |     |       |
| rrsapes       | 12,184 | 146     |     |       |
| dew-satz-1c   | 11,392 | 124     |     |       |
| g2wsat        | 11,060 | 127     |     |       |
| dew-satz-1a   | 10,848 | 127     |     |       |
| lsatv1-1      | 10,837 | 61      |     |       |
| dew-satz-1b   | 10,543 | 125     |     |       |
| ranov         | 10,360 | 134     |     |       |
| rpaws10       | 9,420  | 99      |     |       |
| wllsatv1      | 7,727  | 86      |     |       |
| saps          | 5,735  | 95      |     |       |

(C) UNSAT

| Solver        | Score  | #Solved | Sat | Unsat |
|---------------|--------|---------|-----|-------|
| tts-3-0       | 55,211 | 78      |     |       |
| satelitegti   | 46,427 | 145     |     |       |
| csat          | 35,227 | 113     |     |       |
| zchaff-rand   | 34,597 | 117     |     |       |
| vallst-sh     | 33,065 | 129     |     |       |
| minisat-static| 31,366 | 136     |     |       |
| march-dl      | 25,253 | 111     |     |       |
| jerusat1-31-a | 24,387 | 109     |     |       |
| hsat-5        | 20,216 | 105     |     |       |
| hsat-1        | 20,056 | 105     |     |       |
| zchaff        | 18,849 | 97      |     |       |
| jerusat1-31-b | 15,385 | 100     |     |       |

| Solver        | Score  | #Solved | Sat | Unsat |
|---------------|--------|---------|-----|-------|
| hsatrr        | 14,637 | 66      |     |       |
| haifasat2     | 14,525 | 96      |     |       |
| haifasat      | 14,427 | 95      |     |       |
| dew-satz-1b   | 12,669 | 67      |     |       |
| dew-satz-1a   | 12,173 | 66      |     |       |
| sat4j.jar     | 11,204 | 73      |     |       |
| compsat       | 10,613 | 83      |     |       |
| dew-satz-1c   | 9,902  | 63      |     |       |
| lsatv1-1      | 8,551  | 36      |     |       |
| wllsatv1      | 6,000  | 50      |     |       |
| kcnfs-2004    | 4,205  | 49      |     |       |
| tts-3-0       | 1,352  | 20      |     |       |
Table 4. First stage results, INDUSTRIAL category

(A) SAT+UNSAT

| Solver      | Score Sat | Score Unsat | #Solved Sat | #Solved Unsat |
|-------------|-----------|-------------|-------------|--------------|
| satelitegti| 85,602    | 117         | 78          |              |
| minisat-static| 55,638  | 114         | 74          |              |
| haifasat   | 39,359    | 99          | 79          |              |
| haifasat2  | 31,312    | 94          | 77          |              |
| zchaff-rand| 30,400    | 91          | 78          |              |
| zchaff     | 27,201    | 90          | 60          |              |
| compsat    | 26,955    | 88          | 67          |              |
| csat       | 25,106    | 94          | 78          |              |
| jerusat1-31-b| 22,358  | 96          | 65          |              |
| jerusat1-31-a| 18,885  | 94          | 65          |              |
| hsat-5     | 18,593    | 91          | 45          |              |
| hsat-1     | 17,645    | 91          | 45          |              |
| sat4j.jar  | 16,793    | 88          | 61          |              |
| vallst-sh  | 15,271    | 80          | 63          |              |
| wllsatv1   | 12,623    | 82          | 5           |              |
| hsattr     | 11,820    | 66          | 29          |              |
| march-dl   | 10,051    | 67          | 39          |              |
| dew-satz-1a| 7,116     | 61          | 24          |              |
| rrsaps     | 6,028     | 71          | 0           |              |
| dew-satz-1c| 5,795     | 60          | 5           |              |
| dew-satz-1b| 5,530     | 60          | 5           |              |
| vw         | 5,369     | 67          | 0           |              |
| g2wsat     | 5,291     | 65          | 0           |              |
| rpaws10    | 5,248     | 65          | 0           |              |
| ranov      | 4,563     | 59          | 0           |              |
| lsatv1-1   | 4,117     | 56          | 0           |              |
| saps       | 3,442     | 56          | 0           |              |
| kcnfs-2004 | 3,178     | 40          | 0           |              |
| adaptnovelty| 3,154   | 50          | 0           |              |
| tts-3-0    | 1,544     | 28          | 0           |              |

(B) SAT

| Solver      | Score Sat | #Solved Sat | Solver | Score Sat | #Solved Sat |
|-------------|-----------|-------------|--------|-----------|-------------|
| satelitegti| 52,497    | 117         | minisat-static| 36,651 | 114         |
| haifasat    | 19,051    | 99          | haifasat2    | 15,821 | 94          |
| compsat     | 17,576    | 88          | zchaff       | 15,350 | 90          |
| jerusat1-31-b| 14,598  | 96          | zchaff-rand  | 14,290 | 91          |
| csat        | 13,416    | 94          | vallst-sh    | 11,605 | 82          |
| jerusat1-31-a| 11,262  | 94          | hsat-5       | 10,739 | 91          |
| hsat-1      | 9,809     | 91          | sat4j.jar    | 8,445  | 88          |
| vallst-sh   | 7,080     | 80          | vallst-sh    | 7,000  | 80          |
| wllsatv1    | 6,028     | 71          | rrsaps       | 6,028  | 71          |
| hsattr      | 5,960     | 66          | hsattr       | 5,960  | 66          |
| march-dl    | 5,369     | 67          | vw           | 5,369  | 67          |
| g2wsat      | 5,291     | 65          | g2wsat       | 5,291  | 65          |
| rpaws10     | 5,248     | 65          | rpaws10      | 5,248  | 65          |
| march-dl    | 5,175     | 67          | march-dl     | 5,175  | 67          |
| dew-satz-1a | 4,647     | 60          | dew-satz-1a  | 4,647  | 60          |
| ranov       | 4,563     | 59          | ranov        | 4,563  | 59          |
| dew-satz-1b | 4,422     | 60          | dew-satz-1a  | 4,422  | 60          |
| saps        | 3,442     | 56          | saps         | 3,442  | 56          |
| kcnfs-2004  | 3,178     | 40          | kcnfs-2004   | 3,178  | 40          |
| adaptnovelty| 3,154     | 50          | adaptnovelty | 3,154  | 50          |
| tts-3-0     | 1,544     | 28          | tts-3-0      | 1,544  | 28          |

(C) UNSAT

| Solver      | Score Sat | #Solved Sat | Solver      | Score Sat | #Solved Sat |
|-------------|-----------|-------------|-------------|-----------|-------------|
| satelitegti| 34,481    | 78          | vallst-sh   | 8,266     | 63          |
| haifasat    | 20,802    | 79          | jerusat1-31-a| 7,480  | 65          |
| minisat-static| 20,029  | 74          | hsat-5      | 7,411     | 45          |
| zchaff-rand | 15,985    | 78          | hsat-1      | 7,393     | 45          |
| haifasat2   | 15,324    | 77          | hsattr      | 5,836     | 29          |
| csat        | 12,684    | 78          | march-dl    | 4,576     | 39          |
| zchaff      | 11,631    | 60          | dew-satz-1a | 2,436    | 24          |
| compsat     | 8,470     | 67          | wllsatv1    | 994       | 5           |
| sat4j.jar   | 7,919     | 61          | dew-satz-1b | 907      | 5           |
| jerusat1-31-b| 8,420   | 65          | dew-satz-1c | 947      | 5           |
Table 5. Second stage results, Random category

(A) SAT+UNSAT

| Solver      | Score  | #Solved |
|-------------|--------|---------|
| SAT         |        |         |
|            | Sat    | Unsat   |
| kcnfs-2004 | 95,075 | 92      |
| march-dl   | 27,141 | 56      |
| dsatz-1a   | 22,940 | 68      |
| wllsatv1   | 16,145 | 59      |
| satelitegti| 10,074 | 46      |
| minisat    | 10,058 | 45      |
| UNSAT       |        |         |
|            |         |         |
|            |         |         |
| kcnfs-2004 | 97,930 | 75      |
| march-dl   | 25,228 | 43      |
| dsatz-1a   | 19,456 | 50      |
| wllsatv1   | 12,902 | 45      |
| minisat    | 7,369  | 33      |
| sateLitegti| 7,335  | 33      |

(B) SAT

| Solver     | Score  | #Solved |
|------------|--------|---------|
| ranov      | 163,903| 209     |
| g2wsat     | 101,286| 178     |
| vw         | 76,002 | 170     |
| adaptNovelty| 21,748| 119     |
| saps       | 15,603 | 104     |
| kcnfs-2004 | 14,604 | 92      |
| dsatz-1a   | 8,943  | 68      |
| march-dl   | 7,444  | 56      |
| wllsatv1   | 7,202  | 59      |
| satelitegti| 5,198  | 46      |
| minisat    | 5,147  | 45      |

(C) UNSAT

| Solver     | Score  | #Solved |
|------------|--------|---------|
| kcnfs-2004 | 97,930 | 75      |
| march-dl   | 25,228 | 43      |
| dsatz-1a   | 19,456 | 50      |
| wllsatv1   | 12,902 | 45      |
| minisat    | 7,369  | 33      |
| sateLitegti| 7,335  | 33      |

Branching Strategy, pages 47–59, the new version of march-dl, one of the strong solvers of the SAT 2005 competition (3 silver and 2 bronze medals in the random and crafted categories).

One of the reasons of the success of the random $k$-SAT model in the SAT community is its strong mathematical basis. Oliver Kullmann, one of the judges of the SAT 2005 competition, was in charge of the design of the random category. He reports in The SAT 2005 Solver Competition on Random instances, pages 61-102, the way he built the set of benchmarks to be used for the competition and analyzes the solvers performances on those benchmarks.

The most successful special track of the SAT competition was the Pseudo Boolean evaluation. Vasco Manquinho and Olivier Roussel report in details the event in The First Evaluation of Pseudo-Boolean Solvers (PB’05), pages 103–143. The report is especially interesting because many solvers were found incorrect during the evaluation for mainly two reasons: there was no common input format before the evaluation so I/O errors were frequent and the arithmetic on the coefficients needed special attention to avoid overflow. As a consequence, the organizers decided to rerun corrected version of those solvers after the SAT conference. The report is presenting results updated in September 2005.

For the third time, a QBF evaluation was organized. Massimo Narizzano et al. detail in The third QBF solvers comparative evaluation, pages 145–164, the way the set of benchmarks used was designed and classified. The behavior of the solvers on those benchmarks is then discussed. The evaluation once again emphasized the difficulty to check the correctness of a QBF solver: 5 solvers out of 13 answered incorrectly this year.
The SAT 2005 Competitions and Evaluations

Table 6. Second stage results, Crafted category

(A) SAT+UNSAT

| Solver   | Score  | #Solved |
|----------|--------|---------|
| vallst.sh | 56,445 | 138 100 |
| satelitegti | 53,128 | 122 126 |
| march-dl | 52,432 | 138 99  |
| minisat  | 43,691 | 122 121 |
| hsat-1   | 39,497 | 130 90  |
| csat     | 38,324 | 113 112 |
| zchaff   | 27,455 | 112 89  |
| zchaff-rand | 24,171 | 107 78  |
| tts-3-0  | 21,298 | 5 54    |
| jerusat-a| 19,632 | 104 77  |

(B) SAT

| Solver   | Score  | #Solved |
|----------|--------|---------|
| vallst.sh | 31,258 | 138     |
| march-dl | 27,656 | 138     |
| hsat-1   | 20,156 | 130     |
| satelitegti | 17,418 | 122     |
| minisat  | 17,210 | 122     |
| csat     | 13,791 | 113     |
| zchaff   | 13,692 | 112     |
| zchaff-rand | 11,431 | 107     |
| jerusat-a| 10,702 | 104     |
| tts      | 475    | 5       |

(C) UNSAT

| Solver   | Score  | #Solved |
|----------|--------|---------|
| satelitegti | 35,639 | 126     |
| minisat  | 26,159 | 121     |
| vallst.sh | 25,532 | 100     |
| march-dl | 25,371 | 99      |
| csat     | 23,878 | 112     |
| tts-3-0  | 20,765 | 54      |
| hsat-1   | 19,936 | 90      |
| zchaff   | 14,359 | 89      |
| zchaff-rand | 12,419 | 78      |
| jerusat-a| 9,275  | 77      |

Hossein Sheini and Karem Sakallah present their new pseudo boolean solver in Pueblo: A Hybrid Pseudo-Boolean SAT Solver, pages 165–189. The main feature of that extension of the MiniSat solver for pseudo boolean constraints is to use both clauses and pseudo boolean constraints during learning and backjumping. The details needed to implement a similar solver on top of Minisat are provided.

Five research notes complete this special volume. First, Olivier Bailleux et al submitted to the PB 05 evaluation a solver also based on a translation of the original problem into SAT. A Translation of Pseudo Boolean Constraints to SAT, page 191–200, completes the previous work on MiniSat+ by a new encoding.

Armin Biere and Carsten Sinz propose in Connected components in Compsat, pages 201–208 a way to deals with connected components in conflict driven clause learning solvers, as implemented in the SAT 2005 competitor Compsat.

Vasco Manquinho and João Marques-Silva describe in On Using Cutting Planes in Pseudo-Boolean Optimization, pages 209–219 the solver bsolo, a branch and bound PB solver in which cutting planes are used both for conflict analysis and to improve lower bounds.

Klas Markström proposes in Hard SAT-instances and locality, pages 221–227, a new class of problems based on Eulerian graphs whose instances are expected to be hard for resolution-based SAT solvers.
Table 7. Second stage results, Industrial category

| Solver          | Score | #Solved | | #Solved |
|-----------------|-------|---------||---------|
| satelitegti     | 99,662| 180     | 87 | 73,506 | 180 |
| minisat         | 69,485| 166     | 84 | minisat | 50,985 | 166 |
| haifasat        | 50,931| 151     | 91 | jerusat-b | 38,625 | 163 |
| zchaff-rand     | 50,515| 132     | 94 | haifasat | 28,428 | 151 |
| jerusat-b       | 47,487| 163     | 80 | zchaff-rand | 24,885 | 132 |
| csat            | 36,526| 140     | 91 | csat    | 21,997 | 140 |
| zchaff          | 31,702| 121     | 76 | zchaff  | 19,236 | 121 |
| compsat         | 25,399| 114     | 75 | compsat | 16,715 | 114 |
| sat4j           | 21,097| 110     | 70 | sat4j   | 12,898 | 110 |
| hsat-5          | 20,995| 99      | 54 | wllsatv1 | 11,390 | 86 |
| vallst.sh       | 16,874| 85      | 69 | hsat-5  | 11,046 | 99 |
| wllsatv1        | 12,467| 86      | 6  | vallst.sh | 7,757 | 85 |

| Solver          | Score | #Solved | | #Solved |
|-----------------|-------|---------||---------|
| satelitegti     | 27,518| 87      |  |  |
| zchaff-rand     | 26,792| 94      |  |  |
| haifaSat        | 23,666| 91      |  |  |
| minisat         | 19,863| 84      |  |  |
| csat            | 15,892| 91      |  |  |
| zchaff          | 13,829| 76      |  |  |
| jerusat-b       | 10,225| 80      |  |  |
| hsat-5          | 10,029| 54      |  |  |
| vallst          | 9,192 | 69      |  |  |
| compsat         | 9,097 | 75      |  |  |
| sat4j           | 8,654 | 70      |  |  |
| wllsatv1        | 1,053 | 6       |  |  |

Emmanuel Zarpas concludes the special volume by Back to SAT05 competition: an a posteriori analysis of solvers performances on industrial benchmarks, pages 229–237, in which the solvers are compared mainly on the various SAT-based Bounded Model Checking benchmarks coming from the IBM Formal Verification Benchmarks Library.

Special Thanks

We would like to thank the numerous reviewers involved in that special volume: Fadi Aloul, Anbulagan, Gilles Audemard, Olivier Bailleux, Peter Barth, Roberto Bayardo, Armin Biere, Alexander Bochkov, Yacine Boufkhad, Franc Brzeg, Gilles Dequen, Heidi Dixon, Nikas Eén, Uwe Egly, Marijun Heule, Eugene Goldberg, Edward A. Hirsch, Holger Hoos, Andreas Kuehlmann, Oliver Kullmann, Chu-Min Li, Inês Lynce, João Marques-Silva, Massimo Nardizano, Alexander Nadel, Guoqiang Pan, Steven Prestwich, Jussi Rintanen, Olivier Roussel, Karem Sakallah, Carsten Sinz, Niklas Sörensson, Armando Tacchella, Hans van Maaren Miroslav Vele, Ke Xu, Yinlin Yu and Riccardo Zecchina.

We also would like to thank the three judges for their active involvement in all the choices made at each step of the contest. Armin Biere was our consultant on industrial benchmarks and solvers. Olivier Kullmann generated all the benchmarks for the random
category. Allen van Gelder, who proposed the new scoring scheme, was particularly active. We also want to thank the LRI and Michal Kouril (EECS) for providing us with a very large amount of CPU time.

Finally, we would like to thank Marijn Heule for his help with all the editing issues and Hans van Maaren, the editor in chief of JSAT, for all its help for making this volume a reality.

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