A new linear method to find the congestion hyperplane in DEA

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Received: 14 September 2018 / Accepted: 30 January 2019 / Published online: 13 February 2019
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Abstract
DEA is the way of evaluating the performance of decision-making units on the basis of the degree of efficiency. Unfortunately, congestion has been known as a technical inefficiency for at least three decades just because of the lack of determining the congestion border. In this article, we have introduced the concept of the congestion hyperplane without considering the efficiency value. This has considerably reduced the calculation, and the congestion border has been determined. In addition, the existence of this hyperplane is ascertained. For this purpose, we determine the BCC-efficient DMUs. The normal vector of the hyperplane is denoted by these values. This hyperplane can pass from any of the BCC-efficient DMUs. Next, we have shown that the previous congestion methods are covered and improved by this hyperplane.

Keywords Data envelopment analysis (DEA) · Congestion · Hyperplane · Linear programming · Decision-making units (DMU)

Introduction
DEA is a nonparametric mathematical method that is used to evaluate the performance of a set of decision-making units. Congestion occurs when one or more input(s) is/are increased to reduce one or more output(s) without making changes in the other inputs and outputs. Grosskopf and Fare [8] proposed an executable form and studied the related DEA models (Model: FGL) to evaluate the performance of the product. Later, Cooper, Tomston, and Trall (CTT) introduced another model for the study of congestion. Cooper [6] compared these two models using numerical examples. Cooper [4] proposed another theory in which the two-step calculation is combined in CTT. Brackets [2] suggested CTT models to examine the relationship between employment and output in order to increase employment in the manufacturing industries of China [7]. Cooper [3] studied the congestion management in Chinese industry and showed how evaluating managerial inefficiency caused an increase in the output size without reducing employment in the textile and automobile industry. The congestion in the Chinese industry was re-examined by Jahanshahloo and Khodabakhshi [9]; the proposed models create the evaluation principles for the DMUs. Tons and Saho [12] proposed a nonparametric process in order to measure the scale of production–congestion [5]. Cooper examined congestion under a stochastic DEA using chance programming. Odeck [11] discussed the impact of congestion on inputs, such as fuel consumption, along with the impact of the number of workers in the bus industry in Norway. In his article [13] such congestion is investigated with constant returns to scale output-oriented models. Nora et al. [10] presented a new method for calculating the amount of congestion.

So far, conceptual congestion has been localized so that each unit should be examined separately, which is very time-consuming and requires solving programming models. But the introduction of the congestion hyperplane is a linear criterion that calculates the congestion of each unit without solving any model. Since efficient units are the basis for introducing this hyperplane, it can always be introduced without using any loops. So the process of finding the hyperplane will be a linear and always convergent approach. There was no way to determine the congestion border until now. This congestion border makes it possible to recognize congestion from technical inefficiency. While congestion border helps to comprehensively identify congestion units and their amounts, previous methods could recognize congestion in evaluating...
congestion units. Sometimes, the amount of congestion in such units has been found to be zero because of nonconsideration of the convexity condition. For solving all these problems and reducing computation, the congestion hyperplane is introduced and hence congestion is not recognized as a technical inefficiency anymore. The congestion amount of congestion units is determined by the distance between points on the right-hand side of the hyperplane (right points) and the hyperplane, while threshold congestion is defined by the distance between left points and the hyperplane.

In this paper, we first introduce the methods of Cooper et al. [3] and Nora et al. [10] in “Background” section. The existence of the congestion hyperplane is ascertained by some theorems in “Proposed method” section. In “Numerical example” section, we propose a congestion hyperplane based on a linear method by using efficient units, which can separate congestion units from other DMUs. This hyperplane improves the model developed by Nora et al. [10], and the results are similar to those derived from the model of Cooper et al. [3]. Different numerical examples are selected out of two articles [3, 10] and then solved with the proposed method. Then, these values are compared with the previous methods. In the end, the results of this study are presented in “Conclusions” section.

**Background**

**Basic concepts**

**BCC model**

If \( T_V = \{(x, y) | x \geq \sum_{i=0}^{n} \lambda_i x_i, y \leq \sum_{j=1}^{n} \lambda_j y_j, \sum_{j=1}^{n} \lambda_j = 1, \lambda_j \geq 0, j = 1, \ldots, n\} \), then to measure the relative efficiency of each DMU in the output oriented, the following model should be solved:

\[
\begin{align*}
\text{Max} & \quad \varphi \\
\text{s.t.} & \quad (x_i, \varphi y_i) \in T_V
\end{align*}
\]

**BCC model (1) is:**

\[
\varphi^* = \text{Max} \varphi + \epsilon \left( \sum_{i=1}^{m} s^+_{i} + \sum_{j=1}^{s} s^-_{j} \right)
\]

\[
\text{s.t.} \quad \sum_{i=1}^{m} x_i \lambda_i - s^+_{i0} = x_{i0}, \quad i = 1, 2, \ldots, m
\]

\[
\sum_{j=1}^{m} y_j \lambda_j - s^-_{j0} = \varphi y_{j0}, \quad r = 1, 2, \ldots, s
\]

\[
\lambda_j = 1
\]

\[
(\lambda_j, s^+_{io}, s^-_{jo}) \geq 0, \quad j = 1, 2, \ldots, n, \quad i = 1, 2, \ldots, m, \quad r = 1, 2, \ldots, s
\]  \( (1) \)

**Hyperplane**

The result of solving the multiplicative models is to find the weights such as \((U^*, V^*)\) for the inputs and outputs of the unit. Given the constraints, we have a multiplicative form \(U^*Y_j - V^*X_j \leq 0, j = 1, 2, \ldots, n\). In other words, the \(n\) hyperplanes are available with 1 gradient vector \((U, V)\).

**Congestion**

There is a congestion if the decrease in one or more inputs is accompanied by an increase in one or more outputs (congestion in the input), without the other inputs and outputs getting worse or vice versa, if the increase in one or more inputs along with a decrease in one or multiple outputs without getting worse in other inputs and outputs.

**Introducing the Cooper’s method**

In this section, the method of Cooper et al. [3] and their definition of congestion will be explained. Consider that there are \(n\) DMUs that they have \(m\) inputs and \(s\) outputs. The \(x_j = (x_{1j}, x_{2j}, \ldots, x_{mj})^T\) and \(y_j = (y_{1j}, y_{2j}, \ldots, y_{sj})^T\) are the input and output vectors of DMU \(j = 1, 2, \ldots, n\), respectively.

First, they solved the output-oriented BCC (Banker, Charnes, Cooper) [1] model and then the efficiency of each DMU was obtained by solving BCC model (1).

In the above method, inefficiency is a necessary condition for the presence of congestion, and they first determined inefficient units with Model (1) and finally presented the following model to calculate the amount of congestion:

\[
\varphi^* = \text{max} \varphi + \epsilon \left( \sum_{i=1}^{m} s^+_{i} - \epsilon \sum_{j=1}^{s} s^-_{j} \right)
\]

\[
\text{s.t.} \quad \sum_{i=1}^{m} x_i \lambda_i + s^+_{i0} = x_{i0}, \quad i = 1, 2, \ldots, m
\]

\[
\sum_{j=1}^{m} y_j \lambda_j - s^-_{j0} = \varphi y_{j0}, \quad r = 1, 2, \ldots, s
\]

\[
\lambda_j = 1
\]

\[
(\lambda_j, s^+_{io}, s^-_{jo}) \geq 0, \quad j = 1, 2, \ldots, n, \quad i = 1, 2, \ldots, m, \quad r = 1, 2, \ldots, s
\]  \( (2) \)

Model (1) is solved for all units to identify inefficient DMUs, but Model (2) is used to determine the congestion of inefficient units. Pay attention to the input slack. In practice, the calculations should be conducted in three stages and three models must be solved to get the optimal solution as well as to determine the amount of congestion.
Cooper et al. [3] expressed the following definition to determine congestion in a DMU.

**Definition 1** A DMU has congestion if and only if in an optimal solution \((q^*, \lambda^*, S^{*+}, S^{-+})\) of Model (2), at least one of the following two conditions is occurred:

1. \(q^* > 1\), and there exists at least one \(s_{i1}^{*-} > 0 (1 \leq i \leq m)\).
2. There should be at least one \(s_{r1}^{*+} > 0 (1 \leq r \leq s)\) and at least one \(s_{i1}^{+-} > 0 (1 \leq i \leq m)\).

**Introducing Nora’s method**

First, they solved Model (1) for all DMUs and denoted the optimal solution \((q^*, \lambda^*, S^{*+}, S^{-+})\) each DMU; then, they defined Set EF as follows:

\[
\text{EF} = \{j|q^* = 1\} \text{ or } \text{EF} = \{\text{BCC efficient units}\}
\]

There exists at least one DMU, says DMU\(1\) in EF set, that has the highest value in its first input component compared to the first input component of the remaining DMUs of Set EF. That is to say:

\[
\exists (K \in \text{EF}) s.t. \forall (j \in \text{EF}) \rightarrow x_{1K} > x_{1j}
\]

They delineated \(x_{1K}\) by \(x_{11}\). For all input components \(i = 1, \ldots, m\), they introduced a DMU in EF that \(i\)th input is higher than those of all other DMUs in the set and shown by \(x_{1i}\), \(i = 1, 2, \ldots, m\) similarly. Then, the defined congestion is as follows:

**Definition 2** ADMU has congestion if and only if, in an optimal solution \((q^*, \lambda^*, S^{*+}, S^{-+})\) of Model (1) for this DMU, at least one of the following two conditions is occurred:

1. \(q^* > 1\), and there exists at least one \(x_{i0}^{*} > x_{i1}^{*} (1 \leq i \leq m)\).
2. There is at least one \(s_{r1}^{*+} > 0 (1 \leq r \leq s)\) and at least one \(x_{i0}^{*} > x_{i1}^{*} (1 \leq i \leq m)\).

where \(x_{i0}^{*} > x_{i1}^{*}\); they show the amount of congestion in the \(i\)th input of DMU\(0\) by \(s_{i1}^{*}\) and it is defined as follows:

\[
s_{i1}^{*} = x_{i0}^{*} - x_{i1}^{*}
\]

A DMU has no congestion when \(x_{i0}^{*} \leq x_{i1}^{*}\) and \(s_{i1}^{*} = 0\).

**Proposed method**

In this section, we prove the existence of the congestion hyperplane. In other words, there is a boundary that divides the possible production set (PPS) units into two sections: congestion units and units without congestion (noncongestion units). These two sets have no intersection. Assume that \(T\) is the set of congestion units defined by Cooper et al. [3] and \(T'\) is the set of noncongestion units, according to this definition.

**Theorem 1** \(T\) and \(T'\) sets are convex, \(T\) is the set of congestion units, and \(T'\) is the set of noncongestion units defined by Cooper et al. [3].

**Proof** We consider two arbitrary units to prove that Set \(T\) is convex:

1. If both these DMUs are efficient.

Now, we show that any convex combination of these two members of Set \(T\) will belong to \(T\). It means:

\[
\forall (x_{11}, y_{11}), (x_{21}, y_{21}) \in T, \ (\lambda x_{11} + (1 - \lambda) x_{21}, \lambda y_{11} + (1 - \lambda) y_{21}) \in T,
\]

suppose that:

\[
\begin{align*}
\text{DMU}_1 &= (x_{11}, y_{11}) = (x_{11}, \ldots, x_{1m}, y_{11}, \ldots, y_{1s}) \\
\text{DMU}_2 &= (x_{21}, y_{21}) = (x_{21}, \ldots, x_{2m}, y_{21}, \ldots, y_{2s})
\end{align*}
\]

According to Definition 1 of congestion, if a DMU is efficient, then there is a component greater than zero in congestion slack and excess slack vectors. Thus, we consider that these components for DMU\(_1\) are \(s_{11}^{-*} > 0\), \(s_{1r}^{*-} > 0\), while the components for DMU\(_2\) are \(s_{22}^{*-} > 0\), \(s_{2r}^{*-} > 0\). Now, we show any convex combination of these two DMUs has congestion. In other words, there is a component greater than zero in congestion slack and the excess vectors corresponding to the convex combination of these two DMUs. The excess vectors \(s_1^{*+}, s_2^{*+}\) are introduced as follows:

\[
\begin{align*}
s_1^{*+} &= (s_{11}^{*+}, \ldots, s_{1r}^{*+}, \ldots, s_{1s}^{*+}) \\
s_2^{*+} &= (s_{21}^{*+}, \ldots, s_{2r}^{*+}, \ldots, s_{2s}^{*+})
\end{align*}
\]

Then, one of these two cases will occur:

(a) If the same components of excess vectors in these DMUs are greater than zero, then \(r_1 = r_2 = r\). Since \(s_{2r}^{*+} > 0\), \(s_{1r}^{*+} > 0\) and \(\lambda > 0\) so we have \(\lambda s_{2r}^{*+} + (1 - \lambda)s_{1r}^{*+} > 0\) and \(\lambda s_{2s}^{*+} + (1 - \lambda)s_{1s}^{*+} > 0\).

(b) If the nonzero components of excess vectors in these DMUs are different, then \(r_1 \neq r_2\):
\[ \lambda s_1^{**} + (1 - \lambda)s_2^{**} = \lambda \left( 0, \ldots, s_{1i}^{**}, 0, \ldots, 0 \right) \\
+ (1 - \lambda) \left( 0, \ldots, s_{2i}^{**}, 0, \ldots, 0 \right) \\
= \left( 0, \ldots, \lambda s_{1i}^{**}, 0, \ldots, (1 - \lambda)s_{2i}^{**}, 0, \ldots, 0 \right) > 0. \]

So if \( s_{2i}^{**} > 0, s_{1i}^{**} > 0 \), then \( \lambda s_{1i}^{**} + (1 - \lambda)s_{2i}^{**} \) should be greater than zero. As the same way, we can prove that if \( s_{1i}^{*e}, s_{2i}^{*e} > 0 \) \( 1 \leq i_1, i_2 \leq m \), then \( \lambda s_{1i}^{*e} + (1 - \lambda)s_{2i}^{*e} \) should be greater than zero. Therefore, if both of these DMUs will be efficient, then \( T \) is convex. So:

1. If both these DMUs are efficient, then they have a nonzero component in their congestion slack vectors. In this case (according to the previous section), the congestion slack of a convex combination of these DMUs has at least a nonzero component and hence this combination will be in set \( T \).

2. If DMU \( j \) is efficient and DMU \( i \) is inefficient, then their convex combination will be inefficient, while the congestion slack and excess slack vectors of DMU \( j \) will have at least a nonzero component like \( s_{1i}^{*e} > 0, s_{2i}^{*e} > 0 \) and the congestion slack vector of DMU \( i \) has \( s_{1i}^{*e} > 0 \). Thus, based on the above proof \( \lambda s_{1i}^{*e} + (1 - \lambda)s_{2i}^{*e} \) will be greater than zero.

So in per case, \( (\lambda x_1 + (1 - \lambda)x_2, \lambda y_1 + (1 - \lambda)y_2) \) has congestion and \( T \) will be convex. Similarly, it can be shown that \( T' \) is convex.

**Theorem 2** \( T \) and \( T' \) sets have no intersection where \( T \) is the set of congestion units and \( T' \) is the set of noncongestion units defined by Cooper et al. [3].

**Proof** Suppose \( T \cap T' \neq \emptyset \). Thus, there is DMU \( j \) as a member of \( T \) and \( T' \). According to the assumption, if DMU \( j \) is an efficient congestion unit, then the congestion slack and excess slack vectors have components greater than zero and this is in contradiction with DMU \( j \in T' \). And also if DMU \( j \) is inefficient, we will reach a contradiction and hence this assumption is false.

**Theorem 3** If \( T \) and \( T' \) are convex and \( T \cap T' = \emptyset \), then there is a hyperplane that divides PPS into these sets.

**Proof** Assume Set \( S \) is the difference between \( T \) and \( T' \) as follows:

We first show the Set \( S \) is convex. According to the definition of the \( S \):

\[
\left\{ (t_1, t_2) \in S \rightarrow (t_1, t_2) = (X_1 - X_2, Y_1 - Y_2) \rightarrow \left\{ \begin{array}{l} (X_1, Y_1) \in T \\ (X_2, Y_2) \in T' \end{array} \right. \right. \\
\left\{ (z_1, z_2) \in S \rightarrow (z_1, z_2) = (X_3 - X_4, Y_3 - Y_4) \rightarrow \left\{ \begin{array}{l} (X_1, Y_1) \in T \\ (X_4, Y_4) \in T' \end{array} \right. \right. \\

So:

\[
(\lambda t_1 + (1 - \lambda)z_1, \lambda t_2 + (1 - \lambda)z_2) = \left( \lambda (X_1 - X_2) + (1 - \lambda)(X_3 - X_4), \right. \\
\left. \lambda (Y_1 - Y_2) + (1 - \lambda)(Y_3 - Y_4) \right). \\
\]

According to Theorem 1, \( T \) and \( T' \) are convex so every combination of their members will be in Set \( S \). Namely,

\[
\left\{ \begin{array}{l} \lambda (X_1 - Y_1) + (1 - \lambda)(X_2 - Y_2) \in T \\ \lambda (X_2 - Y_2) + (1 - \lambda)(X_4 - Y_4) \in T' \end{array} \right. \\
(\lambda t_1 + (1 - \lambda)z_1, \lambda t_2 + (1 - \lambda)z_2) \in S \text{ and } S \text{ is a convex set and } (0, 0) \notin S \text{ because of } T \cap T' = \emptyset. \text{ Therefore, there is a hyperplane relying on a convex set:}
\]

\[
\exists 0 \neq P = (P_1', P_2') : (P_1', P_2')(X, Y) = 0
\]

Therefore, the existence of the congestion hyperplane was proven.

**The presentation of congestion hyperplane**

In this section, we introduce the congestion hyperplane. For this purpose, we determine the BCC-efficient DMUs and then define the maximum value of the input components among these DMUs. The normal vector of the hyperplane is denoted by these values. This hyperplane can pass from any of the BCC-efficient DMUs.

First we solve Model (1) for \( \text{DMU}_j \) \( j = 1, 2, \ldots, n \) and determine the optimal solution \( (\varphi^*, \lambda^*, S^{**}, S^{*e}) \). Then, Set \( EF \) is defined as follows:

\[
EF = \{ j | \varphi^* = 1 \} \text{ or } EF = \{ \text{BCC efficient units} \}
\]

There is a DMU in Set \( EF \) (that is called \( \text{DMUI} \)), and this DMU in the first component has the highest amount than other efficient DMUs in Set \( EF \). It is shown as follows:

\[
\exists (l \in EF) : \forall (j \in EF) : x_{1l} > x_{1j}
\]

Then, we determine \( x_{1l} \) with \( x_{1l}^* \). Now we solve Model (3) to find all \( x_{1i}^* \), \( i = 1, 2, \ldots, m \). We introduce \( E^* \) as follows:

\[
\forall 1 \leq i \leq m, \exists (l \in EF) : \forall (j \in EF) : x_{ij} > x_{ij} \rightarrow l \in E^*
\]
DMU$_k \in EF^*$ has the highest amount in the $k$th input component in Set EF and its input and output vectors are $X^*_k = (x^*_{1k}, x^*_{2k}, \ldots, x^*_{1m}, x^*_{2m})$ and $Y_k = (y_{1k}, y_{2k}, \ldots, y_{sk})$, respectively. There is a hyperplane parallel to axis outputs that passes through DMU$_k$. The normal vector of this hyperplane will be obtained by solving the following system of $m + 2$ equations and $m$ unknowns. Consider that the input vector of DMU with the highest value of a first input component (DMU$_i^*$) is $X^*_1 = (x^*_{11}, x^*_{12}, \ldots, x^*_{1m})$, while $X^*_2 = (x^*_{21}, x^*_{22}, \ldots, x^*_{2m})$ is input vector of a DMU with the highest value in the second input component DMU$_i^*$. Similarly, $X^*_m = (x^*_{m1}, x^*_{m2}, \ldots, x^*_{mk})$ is the input vector of a DMU with the highest value of the $m$th input’s component, that is called DMU$_m^*$.

$$
\begin{align*}
N(X^*_1 - X^*_2) &= 0 \\
N(X^*_2 - X^*_3) &= 0 \\
& \vdots \\
N(X^*_{k-1} - X^*_k) &= 0 \\
N(X^*_m - X^*_k) &= 0 \\
N &= 1 \\
N &= > 0
\end{align*}
$$

(5)

Suppose that the hyperplane passes through DMU$_o^*$ \( \in EF^* \). By using the values obtained from Model (5), this congestion hyperplane is:

$$H = \left\{ \left( \begin{array}{c} x \\ y \end{array} \right) | n_1(x_{ij} - x_{ik}) + \cdots + n_m(x_{mj} - x_{mk}) + 0(y_{ij} - y_{1k}) + \cdots + 0(y_{sj} - y_{sk}) = 0 \right\}$$

(6)

In other words, hyperplane equation will be as follows:

$$H = \left\{ \left( \begin{array}{c} x \\ y \end{array} \right) | \sum_{i=1}^{m} n_i(x_{ij} - x_{ik}) + n_k(x_{kj} - x^*_k) + \sum_{j=1}^{s} 0(y_{ij} - y_{sk}) = 0 \right\}$$

(7)

Now, with a proposition we show a congestion unit based on Definition 1 (Cooper et al. [3]) will satisfy Constraint (8). In fact, a congestion unit at the method of Cooper et al. [3] will have congestion with the proposed model.

**Theorem 4** If the following condition is satisfied, then DMU$_j (j = 1, 2, \ldots, n)$ has congestion:

$$\sum_{i=1}^{m} n_i(x_{ij} - x_{ik}) + n_k(x_{kj} - x^*_k) > 0$$

(8)

**Proof** We consider DMU$_o$ as a congestion unit. Suppose that the $k$th congestion slack component of DMU$_o$ is greater than zero ($x^*_k > 0$). Hence, based on the method of Cooper et al. [3], there is a Constraint (a) as well:

$$\sum_{j=1}^{m} \lambda^*_j x_{kj} < x_{ko}$$

(a)

Nora et al. [10] proved that $x^*_i = \sum_{j=1}^{m} \lambda^*_j x_{ij}, i = 1, \ldots, m$. So we have:

$$x_{ko} - x^*_k > \sum_{j=1}^{m} \lambda^*_j x_{kj} - \sum_{j=1}^{m} \lambda^*_j x_{kj} = 0 \rightarrow x_{ko} - x^*_k > 0, \ i = k.$$  

(b)

$$x_{io} - x^*_i \geq \sum_{j=1}^{m} \lambda^*_j x_{ij} - \sum_{j=1}^{m} \lambda^*_j x_{ij} = 0 \rightarrow x_{io} - x^*_i \geq 0, \ i = 1, 2, \ldots, m, \ i \neq k.$$  

(c)

The hyperplane normal vector is greater than zero, so all of its components are positive. As a result, we have:

$$\sum_{i=1}^{m} n_i(x_{io} - x_{ik}) + n_k(x_{ko} - x^*_k) > 0$$

$$i = 1, i \neq k$$

Note that if the hyperplane passes from each efficient units of EF*, the distance of DMU from the congestion hyperplane will not be changed.

**Definition 3** The congestion amount on each DMU is the distance from the congestion hyperplane to this DMU. In other words, if the congestion’s amount of DMU$_j$ in $i$th input component is shown as $s^*_j$, then $x_{ij} - s^*_j$ must be on the congestion hyperplane.
If all $X_i^* (i = 1, 2, \ldots, m)$ should not be exclusive, then we select among them $m$ different DMUs. If there are not $m$ different $\text{DMU}_i^* (i = 1, 2, \ldots, m)$, then there exist more than one congestion hyperplane.

**Numerical example**

In this section, the proposed model will test with selected numerical examples that had been studied by other papers.

**Explaining example**

(a) Consider eight DMUs with one input and one output as A, B, C, D, E, F, G and H. These DMUs are shown in Fig. 1. The horizontal axis identifies the input and the vertical axis of the output.

(b) This example was solved by Cooper et al. [3] using Model (2), the results of which are provided in Table 1.

We now apply our proposed method to solve the same problem. Considering the efficiency of DMUs, we have $\text{Set } \text{EF} = \{A, B, C\}$, where (2) becomes:

\[
X^* = X_c = 3 > x_j (\forall j : (j \in \text{EF})) \rightarrow c \in \text{EF}^*
\]

According to the proposed model, the congestion hyperplane is $X_j - X^* = 0$. We examine the condition of each DMU to the congestion hyperplane:

\[
\begin{align*}
A & : 1 - 3 = -2 < 0 \\
B & : 2 - 3 = -1 < 0 \\
C & : 3 - 3 = 0 \\
D & : 5 - 3 = 2 > 0 \\
E & : 4 - 3 = 1 > 0 \\
F & : 4 - 3 = 1 > 0 \\
G & : 4.5 - 3 = 1.5 > 0 \\
H & : 3 - 3 = 0
\end{align*}
\]

(c) Consider the six DMUs as A, B, C, D, G and R, in Fig. 2. Each DMU has two inputs and one output. The value of output for R is $y = 10$, and the output of other DMUs are:

| DMU | A | B | C | D | E | F | G | H |
|-----|---|---|---|---|---|---|---|---|
| $\varphi^*$ | 1 | 1 | 1 | 2 | 2 | 1.667 | 1.667 | 2 |
| $S'^*$ | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| $S'^{+}$ | 0 | 0 | 0 | 2 | 1 | 1 | 1.5 | 0 |

**Table 1** The results of the Cooper’s method [3]

**Table 2** The results of the proposed model

| DMU | A | B | C | D | E | F | G | H |
|-----|---|---|---|---|---|---|---|---|
| $s'^{inc}$ | 0 | 0 | 0 | 2 | 1 | 1 | 1.5 | 0 |
DMUs is $y = 1$. The horizontal axis identifies the first input and the vertical axis of the second input.

The efficiency and congestion of these DMUs in accordance with the model of Cooper et al. [3] are shown in Table 3. According to this table, $EF = \{A, B, R\}$ and with Model (3):

However, all $x^*_i$, $i = 1, 2, \ldots, m$ are not exclusive so $x^*_1 = x_{1R}$ and $x^*_2 = x_{2B}$. Therefore, the normal vector is $n_1 = \frac{1}{2}, n_2 = \frac{1}{2}$. According to Model (6), DMUR has the highest amount of output. Therefore, the congestion hyperplane will be passing through $R$ point. The equation of congestion hyperplane will be as follows:

$$\frac{1}{2}(x_{ij} - 5) + \frac{1}{2}(x_{2i} - 5) = 0$$

The situation of all DMUs to the hyperplane is:

$$\begin{align*}
A : & -2 < 0 \\
B : & -2 < 0 \\
C : & \frac{5}{2} > 0 \\
D : & \frac{5}{2} > 0 \\
G : & \frac{5}{2} > 0 \\
R : & 0
\end{align*}$$

As can be seen, the points $C$, $D$ and $G$ have congestion that corresponds to the results of the model developed by Cooper et al. [3]. This result is presented in Table 3.

The congestion value of Definition 3 in these DMUs states is given in Table 4:

As can be seen, the results are similar to Cooper et al. [3].

### Applied example

In Table 5, data set for the Chinese automobile and textile industries during the period 1981–1997 is listed that is used by Cooper et al. [3]. The outputs and inputs are defined as follows:

$Y$ is the output measured in units of one million Renminbi, in 1991 prices; $K$ is capital calculated in units of one million Renminbi, also in 1991 prices; and $L$ is labor measured in units of 1000 persons.

All results obtained by using Model (1) are shown in Table 6. Now we solve this example by using our proposed method. For the textile industry, we have:
By comparing the efficiency score of DMUs in the textile industry, we find that DMU 11 and DMU 14 have the maximum input at the first and second components, respectively. Thus, \( x^*_1 = x_{1,11} = 756 \) and \( x^*_2 = x_{2,14} = 25.45 \).

For the automobile industry, we have:

\[
x_1^* = x_{1,17} = 197.8 \quad \text{and} \quad x_2^* = x_{2,17} = 253.96
\]

The results obtained from the proposed model and Nora et al. [10] and the congestion amounts obtained based on Definition 3, \( x'' \), are shown in Table 7. As can be seen, the distance of congestion DMUs from the hyperplane is
positive, while this amount is less than or equal to zero for efficient or technically inefficient DMUs. The distance between DMU and the hyperplane with $P$ value will be shown in the following table.

**Discussions**

The comparison between the methods of Cooper et al. [3] and Nora et al. [10] shows the many advantages of the proposed model that should be discussed individually. According to Definition 1, congestion will occur, while increases in one or more input(s) would cause reduction in one or more output(s) without making changes in the other inputs and outputs. In other words, if we have one input and one output, the congestion can be achieved, while increases in this input cannot gain more output. Consider Unit 2 and Unit 13 in the automotive industry (Table 5). In this example, the congestion hyperplane is defined in passing from DMU14. Based on Table 3, it is clear that the input and output vectors of DMU2 are smaller than the input and output vectors of DMU14. In fact, DMU14 has been able to obtain more output by consuming more input than DMU2. Therefore, any decrease in the output of Unit 2 depends on inefficiency rather than congestion. In other words, by comparing the inputs of Unit 2 (21.16, 412.3) to those of Unit 14, we have $(21.16, 412.3) < (691, 25.45)$. In addition, comparing DMU14 with DMU2 shows that $\epsilon_{DMU_{14}} = 4949.93 > 866.85 = O_{DMU_{2}}$. Similarly, in Unit 13, $I_{DMU_{13}} = (684, 25.08) < (691, 25.45) = I_{DMU_{14}}$ and $O_{DMU_{13}} = 3250.74 < 4949.93 = O_{DMU_{14}}$

The congestion observed in Units 2 and 13, based on the model of Cooper et al. [3], is the type of inefficiency that is wrongly considered as congestion. Owing to Table 7, the results in the textile industry, and the model of Cooper et al. [3] model, DMU15 has congestion. And also, based on the model of Nora et al. [10], this unit has no congestion. But it is shown that the input vector of Unit 15 has a larger component rather than the input vector of Unit 17. Notice that the congestion hyperplane is considered passing from Unit 17. It means that the congestion of this unit cannot be considered as technically inefficient. The model of Nora et al. [10] has considered zero congestion amount for a congestion unit (Unit 15) by ignoring the convexity condition. The model of Cooper et al. [3] introduces 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14 units as congestion units. But by comparing the input and output vectors of these units with those of Unit 17, we find that the input and output components of these units are less than those of Unit 17. Thus, based on the model of Cooper et al. [3], the inefficiency of these units is considered wrongly as the congestion. The congestion is known as a type of inefficiency, but the model of Cooper et al. [3] measured the amount of congestion in four efficient units in such a way that the proposed model solved this problem. Therefore, in some cases, the proposed model corrects the models of Nora et al. [10] and Cooper et al. [3].

**Conclusions**

In this paper, we have shown the existence of the congestion hyperplane by expressing some propositions. This hyperplane can make a complete separation between congestion units and efficient and technically inefficient units. Also, the normal vector has been introduced by using BCC-efficient
units and the highest consumption input components. The congestion hyperplane not only confirms the previous models, but in some cases corrects the models incapacity to introduce the congestion border, inability to distinguish among the congestion units of technical inefficient units and disregard of the congestion units due to the lack of consideration of the convexity condition. The advantages of this proposed method are the determination of congestion units, estimation of the congestion amount by using the distance of the congestion border per unit and the specification of congestion units based on the convexity condition and considerable reduction of computation. Future research could study the creation of the hyperplane on fuzzy data and integer data.

Compliance with ethical standards

Conflict of interest The authors declare that they have no conflict of interest.

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