Type-Preserving Matrices and Security of Block Ciphers

Riccardo Aragona* 1 and Alessio Meneghetti2

1DISIM, Università degli Studi dell’Aquila
Via Vetoio, 67100 Coppito (AQ), Italy
2Dipartimento di Matematica, Università degli Studi di Trento
Via Sommarive 14, 38123 Povo (TN), Italy

Abstract

We provide a new property, called Non-Type-Preserving, for a mixing layer which guarantees protection against algebraic attacks based on the imprimitivity of the group generated by the round functions. Our main result is to present necessary and sufficient conditions on the structure of the binary matrix associated to the mixing layer, so that it has this property. Then we show how several families of linear maps are Non-Type-Preserving, including the mixing layers of AES, GOST and PRESENT. Finally we prove that the group generated by the round functions of an SPN cipher with addition modulo $2^n$ as key mixing function is primitive if its mixing layer satisfies this property. Moreover we generalise the definition of a GOST-like cipher using a Non-Type-Preserving matrix as mixing layer and we show, under the only assumption of invertibility of the S-Boxes, that the corresponding group is primitive.

Keywords: Cryptosystems, mixing layer, group generated by the round functions, primitive groups

MSC 2010: 20B15, 20B35, 94A60

Email addresses: ric.aragona@gmail.com (R. Aragona), alessio.meneghetti@unitn.it (A. Meneghetti)

*The first author is member of of INdAM-GNSAGA (Italy) and he thankfully acknowledges support by DISIM of the University of L’Aquila and by MIUR-Italy via PRIN 2015TW9LSR “Group theory and applications”
1 Introduction

Most modern block ciphers are iterated block ciphers, i.e. are obtained as composition of round functions, and belong to two families of cryptosystems, i.e. Substitution Permutation Networks (SPN) and Feistel Networks (FN). Within each round three permutations of the plaintext space operates, i.e. a non-linear layer and a linear layer which respectively perform confusion and diffusion (see [25]) and a key mixing function which combines the message with the corresponding round key. Most SPNs use the XOR as key mixing, but in many Feistel Networks (e.g. MARS [8], GOST [16], RC6 [23], SEA [26]) and in other block ciphers not belonging to these two families (e.g. IDEA [20]) the key mixing function is the addition modulo $2^n$, for some integer $n$.

Using addition modulo $2^n$ as key mixing function may increase the nonlinearity of a round function. Intuitively, one could take that adding an extra nonlinear layer increases the complexity of attacks. Actually in [21] the authors prove from a theoretical point of view that adopting a key mixing defined by an addition modulo $2^n$ can help to prevent linear cryptanalysis. Then they consider two toy SPNs with the same structure but with key mixing respectively defined by XOR and addition modulo $2^n$ and check from an experimental point of view that the first one is weaker than the latter against linear cryptanalysis. On the other hand in [19] the authors investigate how the use of addition modulo $2^n$ in round functions influences algebraic attacks. Also in [19] statistical and algebraic properties of addition modulo a power of two are studied from a cryptographic point of view.

In this paper, we aim to investigate which properties of the mixing layer are useful to avoid particular classes of algebraic vulnerabilities on a SPN which use addition modulo $2^n$ as key mixing function. Some algebraic properties of the round functions can indeed hide some weaknesses of the corresponding cipher. Firstly, in 1975 Coppersmith and Grossman [13] defined a family of functions which can be used as round functions of a block cipher and studied the permutation group generated by those. Then it has been found out that some group-theoretical properties can reveal weaknesses of the cipher itself. For example, if such group is too small, then the cipher is vulnerable to birthday-paradox attacks (see [18]). Recently, in [9] the authors proved that if such group is of affine type, then it is possible to embed a dangerous trapdoor on it. More relevant, in [22] Paterson built a DES-like cipher whose encryption functions generate an imprimitive group and showed how the knowledge of this trapdoor can be turned into an efficient attack to the cipher. For this reason, a branch of research in Symmetric Cryptography is focused on showing that the group generated by the encryption functions of a given cipher is primitive (see [3, 4, 5, 11, 12, 24, 27, 29]).

Our aim is to provide some classes of mixing layers which guarantee protection against algebraic attacks based on the imprimitivity of the group generated by the round functions. We do so by identifying a necessary and sufficient property on the structure of the binary matrix associated to the mixing layer, and thus proving that several known families of linear maps satisfy it. In particular, we give the definition of Type-Preserving Matrix and we prove that the group generated by the round...
functions of an SPN cipher with invertible S-Boxes and addition modulo $2^n$ as key mixing function is primitive if its mixing layer is not Type-Preserving. Finally we extend the definition of a GOST-like cipher given in [5] using a not Type-Preserving mixing layer instead of a rotation and we prove that the group generated by the round functions is primitive if the S-Boxes are invertible.

The paper is organised as follows. In Section 2 we give our notation, as well as some basic definitions and results concerning block ciphers and primitive permutations groups. In Section 3 we present a new property for mixing layer, called Non-Type-Preserving, and useful to avoid imprimitive attacks on the corresponding block cipher. Then, after having proved our result regarding the necessary and sufficient conditions for a mixing layer to be Non-Type-Preserving, we show how some known mixing layer, such as the case of GOST [16], PRESENT [7] and AES [14], is Non-Type-Preserving. In Section 4 we prove that an SPN which uses addition modulo $2^n$ as key mixing function and a Non-Type-Preserving matrix as mixing layer is primitive. Finally, we generalise a GOST-like cipher using invertible S-Boxes and Non-Type-Preserving mixing layer and we prove its primity.

2 Notation and preliminary results

Let $n \in \mathbb{N}$ and let the plaintext space be $V = \mathbb{F}_2^n$. Let $\text{Sym}(V)$ be the symmetric group acting on $V$, i.e. the group of all permutations on $V$, and by $\text{AGL}(V)$ the group of all affine permutations of $V$, which is a primitive maximal subgroup of $\text{Sym}(V)$.

A block cipher $C$ is a family of key-dependent permutations of $V$

$$\{\varepsilon_K : V \to V, K \in \mathcal{K}\} \in \text{Sym}(V),$$

where $\mathcal{K}$ is the key space, and $|V| \leq |\mathcal{K}|$. The permutation $\varepsilon_K$ is called the encryption function induced by the master key $K$. The block cipher $C$ is called an iterated block cipher if there exists $r \in \mathbb{N}$ such that for each $K \in \mathcal{K}$ the encryption function $\varepsilon_K$ is the composition of $r$ round functions, i.e. $\varepsilon_K = \varepsilon_{1,K} \varepsilon_{2,K} \ldots \varepsilon_{r,K}$. Each round function is a permutation of $V$ depending on a round key derived by a public procedure known as key-schedule from the master key.

Most modern iterated block ciphers belong to two families of cryptosystems: Substitution Permutation Networks, briefly SPN (see e.g. SERPENT [1], PRESENT [7], AES [14]) and Feistel Networks, briefly FN (see e.g. Camelia [2], GOST [16]). In this paper we mainly deal with ciphers of SPN type and we define a class of round functions for iterated block ciphers which is large enough to include the round functions of classical SPN’s.

Let $V = V_1 \times V_2 \times \ldots \times V_\delta$ where, for $1 \leq j \leq \delta$, $\dim(V_j) = m$ and $\times$ represents the Cartesian product of vector spaces. The spaces $V_j$’s are called bricks.

Definition 2.1. For each $k \in V$, a round function induced by $k$ is a map $\varepsilon_k \in \text{Sym}(V)$ where $\varepsilon_k = \gamma \lambda \sigma_k$ and
• $\gamma : V \to V$ is a non-linear permutation, called parallel S-Box, which acts in parallel way on each $V_j$, i.e.

$$(x_1, x_2, \ldots, x_n) \gamma = ((x_1, \ldots, x_m)\gamma_1, \ldots, (x_{m(\delta-1)+1}, \ldots, x_n)\gamma_\delta).$$

The maps $\gamma_j : V_j \to V_j$ are traditionally called S-boxes,

• $\lambda \in \text{Sym}(V)$ is a linear map, called mixing layer,

• $\sigma_k : V \to V$ is the key mixing function, that is a permutation of $V$ combining the message with the corresponding round key $k$.

Since studying the role of the key-schedule is out of the scopes of this paper, we can simply suppose that round keys are randomly-generated vectors in $V$.

Usually, the key mixing function of the well-established SPN’s, such as AES, PRESENT, SERPENT, is $\sigma_k : x \mapsto x + k$, where $+$ is the usual bitwise XOR. Notice that SPN’s featuring a XOR-based key addition have been also called Translation-Based ciphers in [12]. In many other ciphers (e.g. MARS [8], GOST [16], IDEA [20], RC6 [23], SEA [26]) the key mixing is the addition modulo $2^m$, for some integer $m$. This kind of key mixing function may be used to increase the nonlinearity of a round function (see for example [21]). In particular, in this work we are interested in SPN’s which combine the message with the key by the addition modulo $2^{\text{dim}(V)}$ (see [19] [21]).

Definition 2.2. We denote by $\text{SPN}_{\text{mod}}$ an SPN operating on the plaintext space $V$ in which the key mixing function is the addition modulo $2^n$, where $n = \text{dim}(V)$.

2.1 Group generated by the round functions and Primitivity

Besides the classical statistical attacks (e.g. differential and linear cryptanalysis), it is proved that also some algebraic attacks can be effective and dangerous. In this paper we focus on a particular attack, described in [22], based on the imprimitivity of the permutation group generated by the round functions of a block cipher.

Let $C = \{\varepsilon_K \mid K \in K\} \subseteq \text{Sym}(V)$ be an $r$-round iterated block cipher. Several researchers have shown in recent years that the group generated by the encryption functions of a block cipher

$$\Gamma(C) \overset{\text{def}}{=} \langle \varepsilon_K \mid K \in K \rangle \leq \text{Sym}(V)$$

can reveal weaknesses of the cipher itself (see for example [9] [18] [22]). However, the study of $\Gamma(C)$ is not an easy issue in general, since it strongly depends on the key-schedule function (for an example of a key-schedule related study, see [15]). Hence the research focuses on the group generated by the round functions

$$\Gamma_\infty(C) \overset{\text{def}}{=} \langle \varepsilon_{h,K} \mid 1 \leq h \leq r, K \in K \rangle,$$

where all the possible round keys for round $h$ are considered as varying $K \in K$. Such group contains $\Gamma(C)$ and allows to ignore the effect of the key-schedule.
In our case $C$ is an $r$-round SPNmod cipher and the $i$-th round function is $\varepsilon_{i,K} = \gamma^i\sigma_K$, where $K_i$ is the $i$-th round key derived by the key schedule

\[
\begin{align*}
K &\rightarrow V^r \\
K &\mapsto (K_1, \ldots, K_r),
\end{align*}
\]

which we suppose surjective w.r.t. any round.

The corresponding group generated by the round functions is

\[
\Gamma_{\infty} \overset{\text{def}}{=} \langle \gamma^i\sigma_{K_i} \mid 1 \leq i \leq r, K \in K \rangle,
\]

Throughout this paper, sometimes we shall denote $\rho = \gamma\lambda$.

Note that we can consider two group structures on $V$. The first operation is the bitwise XOR, which will be denoted by $+$ and which makes $V$ into a vector space over $F_2$.

The second operation, denoted by $\ominus$, is the sum modulo $2^n$. That is, we represent $a, b \in V$ as $a = (a_0, a_1, \ldots, a_{n-1})$, $b = (b_0, b_1, \ldots, b_{n-1})$, with $a_i, b_i \in \{0, 1\}$ integers, and let $a \ominus b = (c_0, c_1, \ldots, c_{n-1})$,

where

\[
(a_0 + a_12 + a_22^2 + \cdots + a_{n-1}2^{n-1}) \equiv (b_0 + b_12 + b_22^2 + \cdots + b_{n-1}2^{n-1}) \equiv c_0 + c_12 + c_22^2 + \cdots + c_{n-1}2^{n-1} \pmod{2^n},
\]

with $c_i \in \{0, 1\}$ integers. (Here $+$ denotes the ordinary sum of integers.) Therefore $V$ under $\ominus$ is the same thing as the group $\mathbb{Z}_{2^n}$ of integers modulo $2^n$, and we will denote it by $(\mathbb{Z}_{2^n}, \ominus)$.

We recall the following elementary fact we will be using repeatedly without further mention.

**Lemma 2.3.** The subgroups of $(\mathbb{Z}_{2^n}, \ominus)$ are linearly ordered; they are the $\langle 2^q \rangle$, for $0 \leq q \leq n$.

Now we prove the first property of the group generated by the round functions of an SPNmod cipher. Let

\[
\mathcal{T} \overset{\text{def}}{=} \mathcal{T}(V) = \{ \sigma_k : v \mapsto v \ominus k \mid k \in V \}
\]

be the group of $\ominus$-translations on $V$. Note that $\mathcal{T}$ transitively acts on $V$.

**Lemma 2.4.**

\[
\Gamma_{\infty} \overset{\text{def}}{=} \{ \mathcal{T}, \rho \}.
\]

In particular, $\Gamma_{\infty}$ acts transitively on $V$.

**Proof.** If we set $k = 0$, then $\rho \sigma_0 = \rho \in \Gamma_{\infty}$, and so $\rho^{-1} \in \Gamma_{\infty}$. Finally for all $k \in V$, we have $\rho^{-1} \rho \sigma_k = \sigma_k \in \Gamma_{\infty}$. \qed
Permutation groups

We recall some basic notions from permutation group theory. Let $G$ be a finite group acting on the set $V$. For each $g \in G$ and $v \in V$ we denote the action of $g$ on $v$ as $vg$. We denote by $vG = \{ vg | v \in G \}$ the orbit of $v \in V$ and by $G_v = \{ g \in G | vg = v \}$ its stabiliser. The group $G$ is said to be transitive on $V$ if for each $v, w \in V$ there exists $g \in G$ such that $vg = w$. A partition $B$ of $V$ is trivial if $B = \{ V \}$ or $B = \{ \{ v \} | v \in V \}$, and $G$-invariant if for any $B \in B$ and $g \in G$ it holds $Bg \in B$. Any non-trivial and $G$-invariant partition $B$ of $V$ is called a block system. In particular any $B \in B$ is called an imprimitivity block. The group $G$ is primitive in its action on $V$ if $G$ is transitive and there exists no block system. Otherwise, the group $G$ is imprimitive in its action on $V$. We remind the following well-known results whose proofs may be found e.g. in [10].

**Lemma 2.5.** A block of imprimitivity is the orbit $vH$ of a proper subgroup $H < G$ that properly contains the stabiliser $G_v$, for some $v \in V$.

**Lemma 2.6.** If $T$ is a transitive subgroup of $G$, then a block system for $G$ is also a block system for $T$.

In our case, $T$ is a transitive subgroup of $\Gamma_\infty$. Since the map $v \mapsto \sigma_v$ is an isomorphism $(V, \oplus) \rightarrow T$, so we have the following:

**Lemma 2.7.** The subgroups of $T$ are of the form

$$\{ \sigma_u : u \in U \},$$

where $U$ is a subgroup of $(V, \oplus)$.

We obtain

**Lemma 2.8.** If $\Gamma_\infty$ acting on $V$ has a block system, then this consists of the cosets of a $\oplus$-subgroup of $V$, that is, it is of the form

$$\{ W \oplus v : v \in V \}$$

where $W$ is a non-trivial, proper subgroup of $(V, \oplus)$.

**Imprimitivity attack**

The cryptanalysts’ interest into the imprimitivity of the group generated by the round functions of a block cipher arise from the study performed in [22], where it is showed how the imprimitivity of the group can be exploited to construct a trapdoor that may be hard to detect. In particular, the author give an example of a DES-like cipher, which can be easily broken since its round functions generate an imprimitive group, but which is resistant to both linear and differential cryptanalysis.
3 Type-Preserving Matrices

In the first part of this section we will recall some preliminary results proved in [5], and to do so we will adopt the same notation introduced therein. We shall denote

- a subset of $\mathbb{F}_2^m$ of cardinality 1 by a white box;
- a subset of $\mathbb{F}_2^m$ of cardinality $1 < t < 2^m$ by a ruled box;
- the full set $\mathbb{F}_2^m$ by a black box.

We will say that a box has white, ruled or black type.

**Definition 3.1.** Let $D$ be a subset of

$$\mathbb{F}_2^n = V_1 \times V_2 \times \cdots \times V_\delta,$$

where each space $V_i$ has dimension $m$. We shall say that $D$ has a type if

$$D = (D \cap V_1) \times (D \cap V_2) \times \cdots \times (D \cap V_\delta).$$

If $D$ has a type, the type of $D$ will be a sequence of $\delta$ white, ruled or black boxes, where the $i$-th box represents the set $D \cap V_i$.

**Remark 3.2** (Remark 4.9 in [5]). According to Lemma 2.3, a subgroup $D$ of $\mathbb{Z}_2^n$ is of the form $\langle 2^q \rangle$, for some $0 \leq q < n$. Hence a subgroup $D = \langle 2^q \rangle$ of $\mathbb{Z}_2^n$ has one of the following two types.

1. When $q \equiv 0 \pmod{m}$, the subgroup has $n_w$ white boxes and $\delta - n_w$ black boxes, where $0 \leq n_w \leq \delta$ such that $q = n_w m$. Note that there are no white boxes when $q = 0$ (the subgroup is the full group $\mathbb{Z}_2^n$), and there are no black boxes when $q = n$ (the subgroup is $\{0\}$).

2. When $q \not\equiv 0 \pmod{m}$, there is a ruled box which is the box containing the $q$-th bit.

By the previous remark, with a slight abuse of notation, we can refer to the type of any subgroup $D$ in $\mathbb{Z}_2^n$ as the triple $(n_w, n_r, n_b)$, where $n_w$, $n_r$ and $n_b$ are respectively the number of white, ruled and black boxes in the type of $D$ with the following restrictions

$$\begin{align*}
    n_w + n_r + n_b &= \delta \\
    0 &\leq n_w \leq \delta \\
    0 &\leq n_r \leq 1 \\
    0 &\leq n_b \leq \delta
\end{align*} \tag{1}$$

In the next Lemma, proved in [5], we consider the behaviour of the modular sum $\oplus$ with respect to types.

**Lemma 3.3.** If $D$ is a subgroup of $\mathbb{Z}_2^n$ and $\nu \in \mathbb{Z}_2^n$, then $D$ and $\nu \oplus D$ have the same type.
In the remaining part of this section we study the diffusion properties of an invertible mixing layer \( \lambda \), namely how the multiplication by a full-rank binary matrix \( \Lambda \) mixes the bricks \( V_1, \ldots, V_\delta \). To do so, we consider \( \Lambda \) to be a \( \delta \times \delta \) block matrix whose blocks are binary square matrices of order \( m \):

\[
\Lambda = \begin{bmatrix}
A_{1,1} & \cdots & A_{1,\delta} \\
\vdots & \ddots & \vdots \\
A_{\delta,1} & \cdots & A_{\delta,\delta}
\end{bmatrix}
\]

We will also use the notation \( \Lambda((i_1,j_1),(i_2,j_2)) \) for the submatrices of \( \Lambda \):

\[
\Lambda((i_1,j_1),(i_2,j_2)) := \begin{bmatrix}
A_{i_1,j_1} & \cdots & A_{i_1,j_2} \\
\vdots & \ddots & \vdots \\
A_{i_2,j_1} & \cdots & A_{i_2,j_2}
\end{bmatrix}
\]

Observe that if \( A_{i,j} = 0 \) whenever \( i \neq j \), i.e. \( \Lambda \) is a diagonal block matrix, then \( \gamma \lambda \) is a parallel map.

Our interest lies in the image of \( D \subseteq F_2^n \) through the mixing layer, thus we will work with the set \( \text{Im}_{\Lambda,D} = \{ v\Lambda : v \in D \} \). In many cases we will need to work with submatrices of \( \Lambda \), and for the sake of simplicity we will write \( \text{Im}_{\Lambda,D} \Lambda((i_1,j_1),(i_2,j_2)) \) to denote the restriction of the image \( \text{Im}(\Lambda((i_1,j_1),(i_2,j_2))) \) to the set obtained by projecting \( D \) on the coordinates corresponding to the boxes \( j_1, \ldots, j_2 \).

We will study which properties of \( \Lambda \) imply

\[
\text{type}(\text{Im}_{\Lambda,D}) = \text{type}(D). \tag{2}
\]

**Definition 3.4.** A matrix \( \Lambda \in \text{GL}(F_2^n) \), or equivalently the corresponding mixing layer \( \lambda \), satisfying Equation (2) for any \( D \subseteq F_2^n \), is called Type-Preserving. Vice versa, if \( \Lambda \) is not Type-Preserving, then we say that it is Non-Type-Preserving.

In this paper we are mainly interested in the subsets \( D \subseteq F_2^n \), such as the subgroups of \( Z_{2^n} \), with type \((n_w,n_r,n_b)\) satisfying Equation (1). Therefore in the remaining part of this section the subsets \( D \subseteq F_2^n \) are all of this kind. Observe that any \( v \in D \) can be written as the concatenation \((v_w|v_r|v_b)\), where the lengths of \( v_w \), \( v_r \) and \( v_b \) are determined by the type of \( D \). In particular, \( v_w \in F_2^{mn_w}, v_r \in F_2^{mn_r} \) and \( v_w \in F_2^{mn_b} \), with the following properties due to the structure of \( D \):

\[
\begin{align*}
\{ (v_w : \exists v = (v_w|v_r|v_b) \in D) \} &= 1 \\
2 \leq |\{ v_r : \exists v = (v_w|v_r|v_b) \in D \}| &\leq 2^{mn_r} - 1 \\
|\{ v_b : \exists v = (v_w|v_r|v_b) \in D \}| &= 2^{mn_b}
\end{align*}
\]

By (1) we have four cases:

1. \( \text{type}(D) = (0,0,\delta) \)
2. \( \text{type}(D) = (6,0,0) \)
3. \( \text{type}(D) = (n_w,0,\delta-n_w) \)
4. type\((D) = (n_w, 1, \delta - n_w - 1)\)

Cases I and II are trivial, namely all invertible linear maps, i.e. all full-rank matrices, preserve these types: type\((D) = (\delta, 0, 0)\) implies type\((\text{Im}_{D}\lambda) = (\delta, 0, 0)\), and type\((D) = (0, 0, b)\) implies type\((\text{Im}_{D}\lambda) = (0, 0, \delta)\). We will focus on the remaining two cases, starting by case II.

**Lemma 3.5.** Let type\((D) = \text{type}(\text{Im}_{D}\lambda) = (n_w, 0, \delta - n_w)\), where \(n_w \in \{1, \ldots, \delta - 1\}\). Then

\[
\Lambda_{(n_w+1,1):(\delta,n_w)} = 0. \tag{3}
\]

**Proof.** We assume type\((D) = (n_w, 0, \delta - n_w)\) and \(\Lambda_{(n_w+1,1):(\delta,n_w)} \neq 0\). We consider two vectors \(v = (v_w|v_b)\) and \(v' = (v'_w|v'_b)\) in \(D\), with \(v \in \ker(\Lambda_{(n_w+1,1):(\delta,n_w)})\) while \(v'\) is outside of it. Observe that the structure of \(D\) implies that \(v_w = v'_w\), and by applying \(\lambda\) to both we obtain \(v\lambda = (v_w\Lambda_{(1,1):(n_w,n_w)}|0)\) and \(v'\lambda = v\lambda + (0|v_b\Lambda_{(n_w+1,1):(\delta,n_w)})\), here in both cases 0 denotes a string of \(n_b\) zeros. Since the two vectors are different, type\((\text{Im}_{D}\lambda) \neq (n_w, 0, \delta - n_w)\), which contradicts the hypotheses of the Lemma.

The above lemma gives us a necessary property on \(\Lambda\) to have a mixing layer which preserves the type \((n_w, 0, \delta - n_w)\). The next result assures us that this is also sufficient.

**Lemma 3.6.** Let \(n_w \in \{1, \ldots, \delta - 1\}\) and \(\Lambda_{(n_w+1,1):(\delta,n_w)} = 0\). Then \(\lambda\) preserves the type \((n_w, 0, \delta - n_w)\).

**Proof.** We construct \(D\) so that its type would be \((n_w, 0, \delta - n_w)\). Then, any vector \(v \in D\) can be written as a concatenation \((v_w|v_b)\), where \(v_w\) is fixed, while

\[
\{v_b : \exists v = (v_w|v_b) \in D\} = \mathbb{F}_2^{\delta-n_w}. \tag{4}
\]

Due to \(\Lambda_{(n_w+1,1):(\delta,n_w)}\) being the zero matrix, the first \(mn_w\) bits of the image of any \(v \in D\) are equal to \(v_w\Lambda_{(1,1):(n_w,n_w)}\), hence the first \(n_w\) boxes of \(\text{Im}_{D}\lambda\) are white. On the other hand, since \(\lambda\) is invertible, \(\Lambda\) has full rank, which can only be possible by assuming that \(\Lambda_{(n_w+1,1):(\delta,n_w)}\) is invertible. By Equation (4), we therefore have \(\{v_b\Lambda_{(n_w+1,1)+(\delta,n_w)} : \exists v = (v_w|v_b) \in D\} = \mathbb{F}_2^{mn_w}\), from which we conclude that type\((\text{Im}_{D}\lambda) = (n_w, 0, \delta - n_w)\).

Note that in Lemma 3.5 and Lemma 3.6 we did not consider the cases \(n_w = 0\) and \(n_w = \delta\), because they respectively correspond to the cases I and II which we have already discussed.

At last, the case III type\((D) = (n_w, 1, \delta - n_w - 1)\).

**Lemma 3.7.** Let both \(D\) and \(\text{Im}_{D}\lambda\) be of type \((n_w, 1, \delta - n_w - 1)\), where \(n_w \in \{1, \ldots, \delta - 2\}\). Then \(\Lambda\) satisfies the following properties:

(a) \(\Lambda_{(n_w+2,1):(\delta,n_w)} = 0\).

(b) \(\Lambda_{(n_w+1,1):(n_w+1,n_w)}\) is not a full-rank matrix.
(c) \( 2 \leq |\text{Im}_{D} (\Lambda_{(n_w+1,n_w+1);(\delta,n_w+1)})| < 2^{m(\delta-n_w-1)} \)

(d) \( |\text{Im}_{D} (\Lambda_{(n_w+1,n_w+2);(\delta,\delta)})| = 2^{m(\delta-n_w-1)} \)

Proof. We proceed similarly to what we did in the Proof of Lemma 3.5, and we divide the proof in four steps, assuming each time that one among Properties (a), (b), (c) and (d) would not be necessary. We use again the notation \( v = (v_w|v_r|v_b) \), where the length of the three vectors depends on the type of \( D \), and we recall that \( v_w \) is the same for each \( v \in D \).

Firstly, we look at what happens if we deny Property (a). In this case, we consider \( v = (v_w|v_r|v_b) \) and \( v' = (v_w|v_r'|v_b') \) in \( D \) with \( v_b \in \ker \Lambda_{(n_w+2,1);(\delta,n_w)} \) and \( v_b' \notin \ker \Lambda_{(n_w+2,1);(\delta,n_w)} \). It follows that the first \( mn_w \) bits of \( v_b \Lambda \) are different from the first \( mn_w \) bits in \( v_b' \Lambda \), hence the first \( n_w \) boxes in \( \text{Im}_{D} \lambda \) are not white, and so the type of \( D \) is not \((n_w,1,\delta-n_w-1)\).

Similarly, if we deny the second property, we have the same conclusion by choosing \( v = (v_w|v_r|v_b) \) and \( v' = (v_w|v_r'|v_b) \), with \( v_r \neq v_r' \).

We do not go through the entire proofs of Properties (c) and (d), since the argument is quite similar to what we already did above. The difference is that we need to use the entire \( D \) instead of just two vectors \( v \) and \( v' \), and therefore prove that \( \text{Im}_{D} \lambda \) does not have respectively a ruled box (by denying Property (c)) and the right number of black boxes (by denying Property (d)).

As we did for Lemma 3.5, we can also prove that the four necessary properties in Lemma 3.7 are also sufficient.

Lemma 3.8. Let \( \Lambda \) be a matrix satisfying the four properties in Lemma 3.7 for a certain integer \( n_w \in \{1, \ldots, \delta-2\} \). Then \( \Lambda \) preserve the type \((n_w,1,\delta-n_w-1)\).

Proof. We consider \( D \) of type \((n_w,1,\delta-n_w-1)\), where its ruled box is the kernel of the matrix \( \Lambda_{(n_w+1,1);(n_w+1,1)} \).

Similarly to what we did for Lemmas 3.5 and 3.6 in Lemma 3.7 and Lemma 3.8 we did not consider the cases \( n_w = 0 \) and \( n_w = \delta-1 \). We discuss these two cases in the following two results.

Lemma 3.9. Let \( \text{type}(D) = \text{type}(\text{Im}_{D} \lambda) = (0,1,\delta-1) \). Then \( \Lambda \) satisfies Properties (e) and (f) of Lemma 3.5, namely

(c) \( 2 \leq |\text{Im}_{D} (\Lambda_{(1,1);(\delta,1)})| < 2^{m(\delta-1)} \)

(d) \( |\text{Im}_{D} (\Lambda_{(1,2);(\delta,\delta)})| = 2^{m(\delta-1)} \)

Conversely, if there exist \( D \) of type \((0,1,\delta-1)\) for which \( \Lambda \) satisfies the two properties above, then \( \Lambda \) preserves the type of \( D \).

Lemma 3.10. Let \( \text{type}(D) = \text{type}(\text{Im}_{D} \lambda) = (\delta-1,1,0) \). Then \( \Lambda \) satisfies Properties (e) and (f) of Lemma 3.5, namely

(b) \( \Lambda_{(\delta,1);(\delta,\delta-1)} \) is not a full-rank matrix.
(c) $\Lambda_{\delta,\delta} \neq 0$.

Conversely, if $\Lambda$ satisfies the two properties above, then there exists $D$ whose type is preserved by $\Lambda$.

Observe that the properties described in Lemmas 3.9 and 3.10 are particular cases of the ones presented in Lemma 3.7. We omit the proofs of the lemmas, since they can be obtained using the same arguments applied to prove Lemma 3.7 and Lemma 3.8. Hence, we denoted the new properties in the same way, and, with a slight abuse of notation, in the following we will simply refer to Lemma 3.7 and its properties, even though when speaking of types $(0,1,\delta - 1)$ and $(\delta - 1,1,0)$ we should be careful and use the dedicated results.

Putting everything together, we obtain our main result, whose proof we omit since it is a straightforward consequence of Lemmas 3.5, 3.6, 3.7, 3.8, 3.9 and 3.10.

**Theorem 3.11.** The mixing layer $\lambda$ is Type-Preserving with respect to the subsets of $F^n_2$ with type $(n_w, n_r, n_b)$ satisfying Equation (1) if and only if there exists an integer $n_w \in \{0, \ldots, \delta\}$ for which either Equation (3) or the four properties in Lemma 3.7 are satisfied.

Even though the proposed characterisation of Type-Preserving matrix gives us a somewhat difficult method to check whether a mixing layer will preserve the type of a certain $D$, we can use Theorem 3.11 to obtain an equivalent formulation which can be used to derive simpler sufficient conditions to assure that a matrix is Non-Type-Preserving.

**Theorem 3.12.** A mixing layer $\lambda$ is Non-Type-Preserving with respect to the subsets of $F^n_2$ with type $(n_w, n_r, n_b)$ satisfying Equation (1) if and only if, for each integer $n_w \in \{0, \ldots, \delta\}$, both Equation (3) and at least one the four properties in Lemma 3.7 are not satisfied.

We remark that many matrices often used to obtain mixing layers are Non-Type-Preserving, simply because they usually do not satisfy Property (a) of Lemma 3.7.

**Lemma 3.13.** If $\Lambda_{(n_w+2,1):(\delta,n_w)} \neq 0$, for any $n_w \in \{1, \ldots, \delta - 2\}$, then $\Lambda$ is Non-Type-Preserving.

**Proof.** By Theorem 3.12 if both Equation (3) and Property (a) are not satisfied, then $\Lambda$ is Non-Type-Preserving. Note that $\Lambda_{(n_w+2,1):(\delta,n_w)}$ is a submatrix of $\Lambda_{(n_w+1,1):(\delta,n_w)}$, so $\Lambda_{(n_w+2,1):(\delta,n_w)} = 0$ implies $\Lambda_{(n_w+1,1):(\delta,n_w)} \neq 0$.

In the next section we show how some known family of mixing layers are Non-Type-Preserving with respect to the subsets of $F^n_2$ with type $(n_w, n_r, n_b)$ satisfying Equation (1).
3.1 Examples of Non-Type-Preserving mixing layers

In this section we characterise some known class of mixing layers by proving whether they are Non-Type-Preserving mixing Layers with respect to the subsets of $\mathbb{F}_2^n$ whose type satisfy Equation (1). With a slight abuse of notation, any of these mixing layers will simply be denoted as Non-Type-Preserving.

Rotation of a GOST-like cipher

In [5], the mixing layer of a GOST-like cipher is defined as the permutation matrix $\Lambda_s$ with $s \in \{m, \ldots, (\delta - 1)m\}$.

Let $\{e_1, \ldots, e_n\}$ be the canonical basis of $\mathbb{F}_2^n$.

**Definition 3.14.** Let $\pi_s \in \text{Sym}(n)$

$$
\pi_s = \begin{pmatrix}
1 & 2 & \ldots & n \\
\pi_s(1) & \pi_s(2) & \ldots & \pi_s(n)
\end{pmatrix}
$$

such that

- for each $1 \leq x \leq n$, $\pi_s(x) = x + s \mod n$ (5)

where $m \leq s \leq (\delta - 1)m$

The permutation binary matrix associated to $\pi$ is the following circulant matrix

$$
\Lambda_s = \left[\begin{array}{c}
e_{\pi_s(1)} \\
\vdots \\
e_{\pi_s(n)}
\end{array}\right]
$$

**Example 3.15.** In the case of GOST, the actual values of the parameters are: $n = 32$, $m = 4$, $\delta = 8$ and $s = 11$. The right rotation by 11 bits of the GOST cipher is the permutation matrix associated to the following permutation of 32 bits.

$$
\pi_{11} = \begin{pmatrix}
1 & 2 & \ldots & 32 \\
12 & 13 & \ldots & 11
\end{pmatrix}.
$$

The mixing layer associated to $\pi_{11}$ is

$$
\Lambda_{\text{GOST}} = \Lambda_{11} = \left[\begin{array}{cc}
0 & \mathbb{I}_{21} \\
\mathbb{I}_{11} & 0
\end{array}\right],
$$

where we denote the $r \times t$ zero matrix by 0 and the $t \times t$ identity matrix by $\mathbb{I}_t$.

**Proposition 3.16.** Let $\Lambda_s$ be a binary circulant permutation matrix associated to the rotation of $s$ bits. Then $\Lambda_s$ is Non-Type-Preserving if and only if $m \leq s \leq m(\delta - 1)$. 


Proof. We write $\Lambda$ as the block matrix

$$
\Lambda = \begin{bmatrix}
0 & \mathbb{1}_{m\delta-s} \\
\mathbb{1}_s & 0 
\end{bmatrix},
$$

where $\mathbb{1}_t$ is the $t \times t$ identity matrix. We will deal with several cases independently, starting by $s = m$.

In this case, for each $n_w \in \{1, \ldots, \delta - 1\}$ we have $\Lambda_{n_w+1,n_w} = \mathbb{1}_m$, hence Equation (4) is never satisfied. Moreover, it follows that also Property (b) is never satisfied. So, the only possibility left is that $\Lambda$ satisfies both Property (c) and Property (d) of Lemma 3.9, so that $\Lambda$ would preserve a certain set $D$ of type $(0,1,\delta-1)$. However, since $\Lambda_{1,2} = \mathbb{1}_m$, it follows that Property (c) cannot be satisfied by a set of such type.

Let now $s$ be strictly larger than $m$. Then, Property (b) is never satisfied, hence we only need to deal with Lemma 3.4, as we did above. Notice that we can still apply the same argument as we did above, and therefore prove that Property (c) cannot be applied.

These two cases together prove that for any $s \in \{m, \ldots, m(\delta-1)\}$ the rotation of $s$ bits is Non-Type-Preserving. We assume now that $s$ is not inside the interval, and prove that $\Lambda$ is a Type-Preserving matrix. Trivially, if $s = 0$ then $\Lambda$ is the identity matrix, which is a Type-Preserving matrix. In the other possible cases, $\Lambda_{(\delta,1):(\delta,\delta-1)}$ is not a full-rank matrix, and $\Lambda_{(\delta,\delta)} \neq 0$. Then, $\Lambda$ satisfies respectively Property (b) and Property (c) of Lemma 3.10 implying that $\Lambda$ is Type-Preserving.

**Corollary 3.17.** The mixing layer of a GOST-like cipher is Non-Type-Preserving.

**Mixing Layer of PRESENT**

The PRESENT mixing layer $\Lambda_{\text{PRESENT}}$ (see [7]) is a permutation matrix in $\text{GL}_{64}(\mathbb{F}_2)$ defined by

$$
\pi(i) = \begin{cases} 
(16(i-1) \mod 63) + 1 & \text{if } 1 \leq i \leq 63 \\
64 & \text{if } i = 64.
\end{cases}
$$

**Lemma 3.18.** The mixing layer of PRESENT is Non-Type-Preserving.

**Proof.** First, recall that in PRESENT we have 16 bricks of dimension 4. Note that

- the bit of value 1 in position $(13,4)$ is contained in the submatrices $(\Lambda_{\text{PRESENT}})_{(3,1):(16,1)}$ and $(\Lambda_{\text{PRESENT}})_{(4,1):(16,2)}$;
- the bit of value 1 in position $(45,12)$ is contained in the submatrices $(\Lambda_{\text{PRESENT}})_{(n_w+2,1):(16,n_w)}$, for each $n_w \in \{3, \ldots, 10\}$;
- the bit of value 1 in position $(61,16)$ is contained in the submatrices $(\Lambda_{\text{PRESENT}})_{(n_w+2,1):(16,n_w)}$, for each $n_w \in \{11, \ldots, 14\}$.

So $(\Lambda_{\text{PRESENT}})_{(n_w+2,1):(16,n_w)} \neq 0$, for each $n_w \in \{1, \ldots, 14\}$ and hence we can apply Lemma 3.13.
MDS matrix

**Definition 3.19.** A matrix over a finite field which has all the minors not equal to zero is called MDS (Maximum Distance Separable).

**Lemma 3.20.** An MDS mixing layer $\Lambda$ over $\mathbb{F}_{2^m}$, with $m > 1$ the dimension of each S-Box, is Non-Type-Preserving.

**Proof.** By definition it follows $\Lambda_{(n_w+2,1);(\delta,n_w)} \neq 0$ for each $n_w \in \{1, \ldots, \delta - 2\}$, hence we can apply Lemma 3.13. \qed

Mixing Layer of an AES-like cipher

Let

$$\text{MixColumns} = \begin{bmatrix} M & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & M \end{bmatrix} \in \text{GL}_4(\mathbb{F}_{2^m}),$$

where $\delta = 2^t$, for some even integer $t$, and we write the matrix as a $2^t \times 2^t$ block matrix with each block being in $\text{GL}_{2^t}(\mathbb{F}_{2^m})$; in particular, 0 is the zero matrix in $\text{GL}_{2^t}(\mathbb{F}_{2^m})$ and $M$ is an MDS matrix in $\text{GL}_{2^t}(\mathbb{F}_{2^m})$.

With the same notation as above, let

$$\text{ShiftRows} = \begin{bmatrix} \mathbf{0}_1 & \mathbf{0}_2 & \cdots & \mathbf{0}_{2^t-1} & \mathbf{0}_{2^t-2} & \mathbf{0}_{2^t-1} \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ \mathbf{0}_2 & \mathbf{0}_3 & \cdots & \mathbf{0}_{2^t-1} & \mathbf{0}_1 \end{bmatrix} \in \text{GL}_8(\mathbb{F}_{2^m})$$

be a circulant block matrix, where $\mathbf{1}_j$ is the identity element of $\mathbb{F}_{2^m}$ in position $(j, j)$ and the zero element of $\mathbb{F}_{2^m}$ everywhere else.

Let us define the following block matrix in $\text{GL}_8(\mathbb{F}_{2^m})$

$$\Lambda = \text{ShiftRows}^T \cdot \text{MixColumns}^T = \begin{bmatrix} \mathbf{1}_1 \cdot M^T & \mathbf{0}_{2^t} \cdot M^T & \cdots & \mathbf{0}_3 \cdot M^T & \mathbf{0}_{2^t} \cdot M^T \\ \mathbf{0}_{2^t} \cdot M^T & \mathbf{1}_2 \cdot M^T & \cdots & \mathbf{0}_4 \cdot M^T & \mathbf{1}_2 \cdot M^T \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ \mathbf{1}_{2^t-1} \cdot M^T & \mathbf{1}_{2^t-2} \cdot M^T & \cdots & \mathbf{0}_2 \cdot M^T & \mathbf{1}_1 \cdot M^T \end{bmatrix}$$

**Example 3.21.** In the case of AES we have 16 bricks of dimension 8, that is, $\delta = 16$ and $m = 8$. Let

$$\text{MixColumns}_{\text{AES}} = \begin{bmatrix} M & 0 & 0 & 0 \\ 0 & M & 0 & 0 \\ 0 & 0 & M & 0 \\ 0 & 0 & 0 & M \end{bmatrix} \in \text{GL}_{16}(\mathbb{F}_{2^8}),$$

where we write it as a $4 \times 4$ block matrix with each block being in $\text{GL}_4(\mathbb{F}_{2^8})$; in particular, 0 is the zero matrix in $\text{GL}_4(\mathbb{F}_{2^8})$ and

$$M = \begin{bmatrix} 2_x & 3_x & 1_x & 1_x \\ 1_x & 2_x & 3_x & 1_x \\ 1_x & 1_x & 2_x & 3_x \\ 3_x & 1_x & 1_x & 2_x \end{bmatrix} \in \text{GL}_4(\mathbb{F}_{2^8})$$

\[14\]
that is, \( 1 \) can apply Lemma 3.13. Finally, we note that the coefficient \( \Lambda \) is Non-Type-Preserving.

**Proposition 3.22.** \( \Lambda \) is Non-Type-Preserving.

**Proof.** Since \( M \) is an MDS matrix, \( \Lambda_{4,1} \neq 0 \). Therefore, \( \Lambda_{(n_w + 2,1);(16,n_w)} \neq 0 \), for each \( n_w \in \{1, \ldots, \delta - 2\} \), so we can apply Lemma 3.13.

**Corollary 3.23.** The mixing layer of AES is Non-Type-Preserving.

**Proof.** The result directly follows from the previous proposition, anyway we make explicit the algebraic computations in the case of the AES cipher. In [14] the authors define the mixing layer of AES using the left matrix action. Since in this paper we using the right action, we have to consider the transpose of \( \Lambda_{AES} \)

\[
(\Lambda_{AES})^T = \begin{bmatrix}
0_1 \cdot M^T & 0_4 \cdot M^T & 0_3 \cdot M^T & 0_2 \cdot M^T \\
0_4 \cdot M^T & 0_1 \cdot M^T & 0_2 \cdot M^T & 0_3 \cdot M^T \\
0_3 \cdot M^T & 0_2 \cdot M^T & 0_1 \cdot M^T & 0_4 \cdot M^T \\
0_2 \cdot M^T & 0_3 \cdot M^T & 0_4 \cdot M^T & 0_1 \cdot M^T \\
\end{bmatrix} \in GL_{16}(\mathbb{F}_{2^8}).
\]

Finally, we note that the coefficient \( (16,1) \) of \( (\Lambda_{AES})^T \) is the coefficient \( (4,1) \) of

\[
\Lambda_4 \cdot M^T = \begin{bmatrix}
0_x & 0_x & 0_x & 0_x \\
0_x & 0_x & 0_x & 0_x \\
0_x & 0_x & 0_x & 0_x \\
1_x & 1_x & 3_x & 2_x \\
\end{bmatrix}
\]

that is, \( 1_x \neq 0_x \). Hence \( (\Lambda_{AES})^T_{(n_w + 2,1);(16,n_w)} \neq 0 \) for each \( n_w \in \{1, \ldots, 14\} \), so we can apply Lemma 3.13.

**4 Applications**

In this section we consider an SPNmod cipher with Non-Type-Preserving mixing layer and we generalise a GOST-like cipher using this kind of mixing layer and, applying the results given in the previous section, we prove that the corresponding group generated by the round functions are both primitive under the only hypothesis on the invertibility of the S-Boxes.
4.1 Primitivity of an SPNmod cipher

In this section we prove that an SPNmod cipher with invertible S-Boxes and Non-Type-Preserving mixing layer is primitive.

Let

\[ V = V_1 \times V_2 \times \ldots \times V_δ = V_1 \parallel V_2 \parallel \ldots \parallel V_δ \]

where \( \parallel \) denotes concatenation of strings and, for \( 1 \leq j \leq δ \), \( \dim(V_j) = m \).

**Theorem 4.1.** Let \( C \) be an SPNmod cipher acting on the plaintext space \( V \), in which a round function has the form

\[ \varepsilon_k = γλσ_k, \]

for the round key \( k \in V \), where

- \( γ \in \text{Sym}(V) \) is a non-linear permutation (parallel S-Box) which acts in parallel way on each \( V_j \), i.e.

\[ (x_1, x_2, \ldots, x_n)γ = (x_1, x_2, \ldots, x_m)γ_1, \ldots, (x_m(δ-1)+1, \ldots, x_n)γ_δ, \]

where \( γ_j \in \text{Sym}(V_j) \).

- \( λ \in \text{Sym}(V) \) is a Non-Type-Preserving mixing layer.

- \( σ_k \in \text{Sym}(V) \) is the \( \oplus \)-translation of \( V \) by \( k \), i.e. \( vσ_k = v \oplus k \), for any \( v \in V \).

Then \( Γ_∞ = Γ_∞(C) \) is primitive.

**Proof.** Recall that \( ρ = γλ \). By Lemma 4.2 we have \( Γ_∞ = \langle T, ρ \rangle \) so, according to Lemma 2.8 in order to prove that \( Γ_∞ \) is primitive, we have to show that there is no non-trivial proper subgroup \( D \) of \( (V, \oplus) \) and \( v \in V \) such that

\[ Dρ = v \oplus D \]

Since \( 0 \in D \), we can take \( v = 0ρ \), hence it is enough to prove that if \( D \neq \{0\} \) is a proper subgroup of \( Z_{2^n} \), then \( Dρ \neq 0ρ \oplus D \). Clearly an invertible parallel S-Box maps any set having a type to another set having the same type, since each S-box is a bijection. Hence \( D \) and \( Dγ \) share the same type and, by Lemma 3.3 this is the same type as \( 0ρ \oplus D \). Therefore \( 0ρ \oplus D \) cannot be equal to \( Dρ \), if we prove that, for any non-trivial proper subgroup \( D \) of \( Z_{2^n} \), \( Dγ \) and \( Dρ = (Dγ)λ \) have different types. Finally, the latter statement follows from Theorem 3.12 since by hypothesis \( λ \) is Non-Type-Preserving.

4.2 Generalised GOST-like cipher and primitivity

In this section, first we extend the definition of GOST-like cipher given in [5] using any Non-Type-Preserving mixing layer instead of a rotation, then we prove that the group generated by the round functions is primitive if the S-Boxes are invertible.
We give the definition of a Generalized GOST-like cipher and of the corresponding group generated by the round functions, arranging the definition of a GOST-like cipher given in [5] by substituting the rotation by \( m \leq s \leq m(\delta - 1) \) with any Non-Type-Preserving mixing layer.

The plaintext space is \( V = V^1 \times V^2 \), where \( V^1, V^2 \) are two copies of \( \mathbb{F}_2^n \), and the key space \( \mathcal{K} \) is another copy of \( \mathbb{F}_2^n \). Clearly \( V \) inherits both group structures componentwise from \( V^1, V^2 \).

Let us consider

- \( V^i \), for \( i = 1, 2 \), as the Cartesian product
  \[
  V^i = V^i_1 \times \cdots \times V^i_\delta = V^i_1\cdots|V^i_\delta
  \]
  of \( \delta > 1 \) spaces \( V^i_j \), all of the same dimension \( m > 1 \).
- a non-linear map (parallel S-Box) \( \gamma \in \text{End}(V^i) \) which acts in parallel way on each \( V^i_j \), where \( \gamma_j \in \text{Sym}(V^i_j) \).
- a Non-Type-Preserving linear map \( \lambda \in \text{Sym}(V^i) \).
- \( \rho \overset{\text{def}}{=} \gamma \lambda \in \text{Sym}(V^i) \).

For \( (k_1, k_2) \in V = V^1 \times V^2 \), consider the \( \oplus \)-translation on \( V \) by \( (k_1, k_2) \)

\[
\sigma_{(k_1, k_2)} : V^1 \times V^2 \rightarrow V^1 \times V^2
\]
\[
(x_1, x_2) \mapsto (x_1 \oplus k_1, x_2 \oplus k_2).
\]

We now introduce a formal \( 2n \times 2n \) matrix, which implements the Feistel structure,

\[
P = \begin{bmatrix}
0 & 1 \\
1 & \rho
\end{bmatrix},
\]

where 0 and 1 are \( n \times n \) matrices. This acts (on the right) on \( (x_1, x_2) \in V = V^1 \times V^2 \) by

\[
(x_1, x_2)P = (x_2, x_1 + x_2\rho).
\]

We are ready to define a round function of a Generalised GOST-like cipher. Let \( \mathcal{H} = \mathcal{K} \times \mathcal{K} = V \) be the key space, a round takes the form

\[
\sigma_k P \sigma_h,
\]

with \( k, h \in \mathcal{H} \).

The corresponding group generated by the round functions will thus be

\[
\Gamma_\infty = \{ \sigma_k P \sigma_h : k, h \in \mathcal{H} \}
\]

**Theorem 4.2.** Let \( \mathcal{C} \) be a Generalised GOST-like cipher as defined above. If the parallel S-Box \( \gamma \) is a permutation of \( V^i \), in other words \( \gamma \in \text{Sym}(V^i) \), then \( \Gamma_\infty(\mathcal{C}) \) is primitive.

**Proof.** The proof is the same as that given in Section 4 of [5], which uses the Goursat’s Lemma [17], until the case \( D\rho = 0\rho \oplus D \) with \( D \) a non-trivial proper subgroup of \( \mathbb{Z}_{2^n} \). Finally, for this case we can proceed as done in the proof of Theorem 4.1 and apply Theorem 3.12.
References

[1] R. J. Anderson, E. Biham, and L. R. Knudsen, SERPENT: A new block cipher proposal, Fast Software Encryption, 222–238, Lecture Notes in Comput. Sci. 1372, Springer, Berlin (1998).

[2] K. Aoki, et al. Camellia: A 128-bit block cipher suitable for multiple platforms-design and analysis, Selected Areas in Cryptography. 39–56, Lecture Notes in Comput. Sci., 2012, Springer, Berlin (2000).

[3] R. Aragona, M. Calderini, A. Tortora, and M. Tota, On the primitivity of PRESENT and other lightweight ciphers, Journal of Algebra and Its Applications, Online Ready (2017), 1–18.

[4] R. Aragona, A. Caranti, F. Dalla Volta, and M. Sala, On the group generated by the round functions of translation based ciphers over arbitrary fields, Finite Fields Appl. 25 (2014), 293–305.

[5] R. Aragona, A. Caranti, and M. Sala, The group generated by the round functions of a GOST-like cipher, Ann. Mat. Pura Appl., 196 (2016), no. 1, 1–17.

[6] A. Bannier, N. Bodin, and E. Filiol, Partition-Based Trapdoor Ciphers, IACR Cryptology ePrint Archive, Report 2016/493 (2016); available at http://eprint.iacr.org/2016/493.

[7] A. Bogdanov et al., PRESENT: An ultra-lightweight block cipher, CHES '07, 450–466, Lecture Notes in Comput. Sci. 4727, Springer, Berlin (2007).

[8] C. Burwick, et al. MARS-a candidate cipher for AES, NIST AES Proposal 268 (1998).

[9] M. Calderini, and M. Sala Elementary abelian regular subgroups as hidden sums for cryptographic trapdoors, preprint, arXiv:1702.00581 [math.GR] (2017).

[10] P. J. Cameron, Permutation groups, London Mathematical Society Student Texts 45, Cambridge University Press, Cambridge (1999).

[11] A. Caranti, F. Dalla Volta, and M. Sala, An application of the O’Nan-Scott theorem to the group generated by the round functions of an AES-like cipher, Des. Codes Cryptogr. 52 (2009), no. 3, 293–301.

[12] A. Caranti, F. Dalla Volta, and M. Sala, On some block ciphers and imprimitive groups, Appl. Algebra Engrg. Comm. Comput. 20 (2009), no. 5-6, 339–350.

[13] D. Coppersmith and E. Grossman, Generators for certain alternating groups with applications to cryptography, SIAM J. Appl. Math. 29 (1975), no. 4, 624–627.

[14] J. Daemen and V. Rijmen, The design of Rijndael: AES – the Advanced Encryption Standard, Information Security and Cryptography, Springer-Verlag, Berlin (2002).
[15] S. M. Dehnavi, A. M. Rishakani, M. M. Shamsabad, H. Maimani, E. Pasha, Cryptographic Properties of Addition Modulo $2^n$. IACR Cryptology ePrint Archive 181 (2016).

[16] V. Dolmatov, GOST 28147-89: encryption, decryption, and message authentication code (MAC) algorithms, Technical report (2010), http://tools.ietf.org/html/rfc5830.

[17] E. Goursat, Sur les substitutions orthogonales et les divisions régulières de l'espace, Ann. Sci. École Norm. Sup. 3(6) (1889), 9–102.

[18] Jr. B. S. Kaliski, R. L. Rivest, and A. T. Sherman, Is the Data Encryption Standard a group? (Results of cycling experiments on DES), J. Cryptology 1 (1988), no. 1, 3–36.

[19] O. Kazymyrov, R. Oliynykov, H. Raddum, Influence of addition modulo $2^n$ on algebraic attacks, Cryptogr. Commun. 8 (2016), no. 2, 277–289.

[20] X. Lai, J. L. Massey, A proposal for a new block encryption standard, Advances in cryptology – EUROCRYPT '90, 389–404, Lecture Notes in Comput. Sci. 473, Springer, Berlin (1990).

[21] D. Mukhopadhyay, D. RoyChowdhury. Key Mixing in Block Ciphers through Addition modulo $2^n$, IACR Cryptology ePrint Archive 383 (2005).

[22] K.G. Paterson, Imprimitive permutation groups and trapdoors in iterated block ciphers, Fast Software Encryption, 201–214, Lecture Notes in Comput. Sci. 1636, Springer, Berlin (1999).

[23] R. L. Rivest, M. J. W. Robshaw, R. Sidney, Y. L. Yin, The RC6TM block cipher. In First Advanced Encryption Standard (AES) Conference (1998).

[24] R. Sparr and R. Wernsdorf, Group theoretic properties of Rijndael-like ciphers, Discrete Appl. Math. 156 (2008), no. 16, 3139–3149.

[25] C. E. Shannon, Communication theory of secrecy systems, Bell System Tech. 28 (1949), 656–715.

[26] F. X. Standaert, G. Piret, N. Gershenfeld, N., J. J. Quisquater, (2006, April). SEA: A scalable encryption algorithm for small embedded applications, Smart Card Research and Advanced Applications – CARDIS '06, 222–236, Lecture Notes in Comput. Sci. 3928, Springer, Berlin, (2006).

[27] R. Wernsdorf, The round functions of RIJNDAEL generate the alternating group, Fast Software Encryption 143-148, Lecture Notes in Comput. Sci. 2365, Springer, Berlin (2002).

[28] R. Wernsdorf, The one-round functions of the DES generate the alternating group, Advances in cryptology-EUROCRYPT’92, Lecture Notes in Comput. Sci. 658, Springer, Berlin (1993).
[29] R. Wernsdorf, *The round functions of SERPENT generate the alternating group* (2000); available at [http://csrc.nist.gov/archive/aes/round2/comments/20000512-rwernsdorf.pdf](http://csrc.nist.gov/archive/aes/round2/comments/20000512-rwernsdorf.pdf)