Improved lower bounds on the number of edges in list critical and online list critical graphs

Hal Kierstead and Landon Rabern*

July 1, 2014

Abstract

We prove that every $k$-list-critical graph ($k \geq 8$) on $n \geq k + 2$ vertices has at least
$$\frac{1}{2} \left( k - 1 + \frac{k^3}{(k-c)(k-1)(k-2)} \right) n$$
edges where $c = (k - 3) \left( \frac{1}{2} - \frac{1}{(k-1)(k-2)} \right)$. This improves the bound established by Kostochka and Stiebitz [15]. The same bound holds for online list-critical graphs, improving the bound established by Riasat and Schauz [20]. Both bounds follow from a more general result stating that either a graph has many edges or it has an Alon-Tarsi orientable induced subgraph satisfying a certain degree condition. We use these lower bounds to prove an Ore-degree version of Brooks’ theorem for online list-coloring: every graph with $\theta \geq 18$ and $\omega \leq \frac{\theta}{2}$ is online $\lfloor \frac{\theta}{2} \rfloor$-choosable. Here $\theta$ is the Ore-degree. Finally, we give an upper bound on the online choice number ($\chi_{OL}$) of $\chi_{OL}$-critical graphs. For an $n$-vertex $d$-regular $\chi_{OL}$-critical graph with independence number $\alpha$ this bound gives $\chi_{OL} < \frac{n}{n-d} + 2$. For triangle-free graphs this bound gives $\chi_{OL} \leq \Delta + 1 - \lfloor \frac{1}{4} \log(\Delta) \rfloor$.

1 Introduction

A $k$-coloring of a graph $G$ is a function $\pi: V(G) \rightarrow [k]$ such that $\pi(x) \neq \pi(y)$ for each $xy \in E(G)$. The least $k$ for which $G$ has a $k$-coloring is the chromatic number $\chi(G)$ of $G$. We say that $G$ is $k$-chromatic when $\chi(G) = k$. A graph is $k$-critical if $\chi(G) = k$ and $\chi(H) < k$ for every proper subgraph $H$ of $G$. If $G$ is $k$-chromatic, then any minimal $k$-chromatic subgraph of $G$ is $k$-critical. In this way, many questions about $k$-chromatic graphs can be reduced to questions about $k$-critical graphs which have more structure. The study of critical graphs was initiated by Dirac [5] in 1951. It is easy to see that a $k$-critical graph $G$ must have minimum degree at least $k - 1$ and hence $2 \|G\| \geq (k - 1) |G|$. The problem of determining the minimum number of edges in a $k$-critical graph has a long history. First, in 1957, Dirac [5] generalized Brooks’ theorem [4] by showing that any $k$-critical graph $G$ with $k \geq 4$ and $|G| \geq k + 2$ must satisfy

$$2 \|G\| \geq (k - 1) |G| + k - 3.$$
In 1963, this bound was improved for large $|G|$ by Gallai [8]. Put

$$g_k(n, c) := \left( k - 1 + \frac{k - 3}{(k - c)(k - 1) + k - 3} \right)n.$$

Gallai showed that every $k$-critical graph $G$ with $k \geq 4$ and $|G| \geq k + 2$ satisfies $2\|G\| \geq g_k(|G|, 0)$. In 1997, Krivelevich [10] improved Gallai’s bound by replacing $g_k(|G|, 0)$ with $g_k(|G|, 2)$. Then, in 2003, Kostochka and Stiebitz [15] improved this by showing that a $k$-critical graph with $k \geq 6$ and $|G| \geq k + 2$ must satisfy $2\|G\| \geq g_k(|G|, (k - 5)\alpha_k)$ where

$$\alpha_k := \frac{1}{2} - \frac{1}{(k - 1)(k - 2)}.$$

In 2012, Kostochka and Yancey [12] drastically improved these bounds by showing that every $k$-critical graph $G$ with $k \geq 4$ must satisfy

$$\|G\| \geq \left\lceil \frac{(k + 1)(k - 2)|G| - k(k - 3)}{2(k - 1)} \right\rceil.$$

Moreover, they show that this bound is tight for $k = 4$ and $n \geq 6$ as well as for infinitely many values of $|G|$ for any $k \geq 5$. This bound has many interesting coloring applications such as a very short proof of Grötzsch’s theorem on the 3-colorability of triangle-free planar graphs [11] and short proofs of the results on coloring with respect to Ore degree in [10, 17, 14]. Given the applications to coloring theory, it makes sense to investigate the same problem as a very short proof of Grötzsch’s theorem on the 3-colorability of triangle-free planar graphs [7]. Let $G$ be a graph and $\pi : V(G) \to \mathbb{N}$ such that $\pi(v) \in L(v)$ for each $v \in V(G)$ and $\pi(x) \neq \pi(y)$ for each $xy \in E(G)$. A graph $G$ is $L$-critical if $G$ is not $L$-colorable, but every proper subgraph $H$ of $G$ is $L|_{V(H)}$-colorable. For $f : V(G) \to \mathbb{N}$, a list assignment $L$ is an $f$-assignment if $|L(v)| = f(v)$ for each $v \in V(G)$. If $f(v) = k$ for all $v \in V(G)$, then we also call an $f$-assignment a $k$-assignment. We say that $G$ is $f$-choosable if $G$ is $L$-colorable for every $f$-assignment $L$. The best known lower bound on the number of edges in an $L$-critical graph where $L$ is a $(k - 1)$-assignment was given by Kostochka and Stiebitz [15] in 2003 and shows that for $k \geq 9$ and $G \not\cong K_k$ an $L$-critical graph where $L$ is a $(k - 1)$-assignment on $G$, we have $2\|G\| \geq g_k(|G|, \frac{1}{3}(k - 4)\alpha_k)$. We improve the bound to $2\|G\| \geq g_k(|G|, (k - 3)\alpha_k)$.

Online list coloring was independently introduced by Zhu [24] and Schaud [21] (Schauz called it paintability). Let $G$ be a graph and $f : V(G) \to \mathbb{N}$. We say that $G$ is online $f$-choosable if $f(v) \geq 1$ for all $v \in V(G)$ and for every $S \subseteq V(G)$ there is an independent set $I \subseteq S$ such that $G - I$ is online $f^*$-choosable where $f^*(v) := f(v)$ for $v \in V(G) - S$ and $f^*(v) := f(v) - 1$ for $v \in S - I$. We say that $G$ is online $f$-critical if $G$ is not online $f$-choosable, but every proper subgraph $H$ of $G$ is online $f|_{V(H)}$-choosable. In 2012, Riasat and Schauz [20] showed that Gallai’s bound $2\|G\| \geq g_k(|G|, 0)$ holds for online $f$-critical graphs where $f(v) := k - 1$ for all $v \in V(G)$. We improve this for $k \geq 8$ by proving the same bound as we have for list coloring: $2\|G\| \geq g_k(|G|, (k - 3)\alpha_k)$.
Our main theorem shows that a graph either has many edges or an induced subgraph which has a certain kind of good orientation. To describe these good orientations we need a few definitions. A subdigraph $H$ of a directed multigraph $D$ is called eulerian if $d_H^-(v) = d_H^+(v)$ for every $v \in V(H)$. We call $H$ even if $\|H\|$ is even and odd otherwise. We write $EE(D)$ (resp. $EO(D)$) for the number of even (resp. odd) spanning subdigraphs of $D$. Note that the edgeless subgraph of $D$ is even and hence we always have $EE(D) > 0$.

Let $G$ be a graph and $f : V(G) \to \mathbb{N}$. We say that $G$ is $f$-Alon-Tarsi (for brevity, $f$-AT) if $G$ has an orientation $D$ where $f(v) \geq d_D^+(v) + 1$ for all $v \in V(D)$ and $EE(D) \neq EO(D)$. Alon and Tarsi \cite{AT} showed that such orientations are very useful for list coloring; they proved the following.

**Lemma 1.1.** If a graph $G$ is $f$-AT for $f : V(G) \to \mathbb{N}$, then $G$ is $f$-choosable.

Schauz \cite{Schauz} extended this result to online $f$-choosability.

**Lemma 1.2.** If a graph $G$ is $f$-AT for $f : V(G) \to \mathbb{N}$, then $G$ is online $f$-choosable.

For a graph $G$, we define $d_0 : V(G) \to \mathbb{N}$ by $d_0(v) := d_G(v)$. The $d_0$-choosable graphs were first characterized by Borodin \cite{Borodin} and independently by Erdős, Rubin and Taylor \cite{ERT}. The connected graphs which are not $d_0$-choosable are precisely the Gallai trees (connected graphs in which every block is complete or an odd cycle). The generalization to a characterization of $d_0$-AT graphs was first given in \cite{Hkat} by Hladký, Král and Schauz.

We prove the following general lemma saying that either a graph has many edges or has an induced $f_H$-AT subgraph $H$ where $f_H$ basically gives the number of colors we would expect the vertices to have left in their lists after $\delta(G)$-coloring $G - H$.

**Theorem 2.13.** Let $G$ be a graph with $\delta := \delta(G) \geq 5$ and $K_{i+1} \not\subseteq G$. If $G$ does not have a nonempty induced subgraph $H$ which is $f_H$-AT where $f_H(v) := \delta + d_H(v) - d_G(v)$ for all $v \in V(H)$, then $2 \|G\| \geq g_{\delta+1}(|G|, c)$ where $c := (\delta - 2)\alpha_{\delta+1}$ when $\delta \geq 7$ and $c := (\delta - 3)\alpha_{\delta+1}$ when $\delta \in \{5, 6\}$.

The Alon-Tarsi number of a graph $AT(G)$ is the least $k$ such that $G$ is $f$-AT where $f(v) := k$ for all $v \in V(G)$. We have $\chi(G) \leq \chi(G) \leq \chi_{OL}(G) \leq AT(G) \leq \text{col}(G)$. We say that $G$ is $k$-AT-critical if $AT(G) = k$ and $AT(H) < k$ for all proper induced subgraphs $H$ of $G$. From Theorem 2.13 we can conclude the following.

**Theorem 3.5.** For $k \geq 6$ and $G \not\supseteq K_k$ a $k$-AT-critical graph, we have $2 \|G\| \geq g_k(|G|, c)$ where $c := (k - 3)\alpha_k$ when $k \geq 8$ and $c := (k - 4)\alpha_k$ when $k \in \{6, 7\}$.

Similarly, applying Lemma 1.1 gives the following.

**Theorem 3.1.** For $k \geq 6$ and $G \not\supseteq K_k$ an $L$-critical graph where $L$ is a $(k - 1)$-assignment on $G$, we have $2 \|G\| \geq g_k(|G|, c)$ where $c := (k - 3)\alpha_k$ when $k \geq 8$ and $c := (k - 4)\alpha_k$ when $k \in \{6, 7\}$.

This improves the bound given by Kostochka and Stiebitz in \cite{KS}; for $L$-critical graphs, they have $2 \|G\| \geq g_k(|G|, \frac{1}{3}(k - 4)\alpha_k)$ for $k \geq 9$. Now, applying Lemma 1.2 gives the following.
Theorem 3.3. For \( k \geq 6 \) and \( G \neq K_k \) an online \( f \)-critical graph where \( f(v) := k - 1 \) for all \( v \in V(G) \), we have \( 2 \|G\| \geq g_k(|G|, c) \) where \( c := (k - 3)\alpha_k \) when \( k \geq 8 \) and \( c := (k - 4)\alpha_k \) when \( k \in \{6, 7\} \).

Definition 1. The Ore-degree of an edge \( xy \) in a graph \( G \) is \( \theta(xy) := d(x) + d(y) \). The Ore-degree of a graph \( G \) is \( \theta(G) := \max_{xy \in E(G)} \theta(xy) \).

A bound like Brooks’ theorem in terms of the Ore-degree was given by Kierstead and Kostochka [10] and subsequently the required lower bound on \( \Delta \) was improved in [17] [14] [18]. For example, we have the following.

Theorem 1.3. Every graph with \( \theta \geq 10 \) and \( \omega \leq \frac{\theta}{2} \) is \( \left\lceil \frac{\theta}{2} \right\rceil \)-colorable.

Another method for achieving the tightest of these results on Ore-degree was given by Kostochka and Yancey [12]. Their proof combined their new lower bound on the number of edges in a color critical graph together with a list coloring lemma derived via the kernel lemma. In Section 4 we improve this latter lemma and, in a similar way, use it in combination with Theorem 3.1 to prove an Ore-degree version of Brooks’ theorem for list coloring. The improved lemma can be seen to be giving another lower bound on the number of edges in \( G \). Let \( \text{mic}(G) \) be the maximum of \( \sum_{v \in I} d_G(v) \) over all independent sets \( I \) of \( G \).

Theorem 4.5. For any graph \( G \) we have either:

1. \( G \) has a nonempty induced subgraph \( H \) which is online \( f_H \)-choosable where \( f_H(v) := \delta(G) + d_H(v) - d_G(v) \) for all \( v \in V(H) \); or
2. \( 2 \|G\| \geq (\delta(G) - 1) |G| + \text{mic}(G) + 1 \).

This quickly gives the aforementioned Ore degree version of Brooks’ theorem for list coloring.

Theorem 5.9. Every graph with \( \theta \geq 18 \) and \( \omega \leq \frac{\theta}{2} \) is \( \left\lceil \frac{\theta}{2} \right\rceil \)-colorable.

Note that using Kostochka and Stiebitz’s above lower bound on the number of edges in a list critical graph gives Theorem 5.9 with \( \theta \geq 54 \). Similarly, we get the online version.

Theorem 5.8. Every graph with \( \theta \geq 18 \) and \( \omega \leq \frac{\theta}{2} \) is online \( \left\lceil \frac{\theta}{2} \right\rceil \)-choosable.

We expect that Theorems 5.8 and 5.9 actually hold for \( \theta \geq 10 \). In the regular coloring case, it was shown in [14] that the only exception when \( \theta \geq 8 \) is the graph \( O_5 \); again, the expectation is that the same result will hold for Theorems 5.8 and 5.9.

Finally, we give a general upper bound on the online choice number that follows immediately from the results in Section 4. The online choice number of a graph \( G \) is the least \( k \) for which \( G \) is online \( f \)-choosable where \( f(v) = k \) for all \( v \in V(G) \). We say that \( G \) is \( \chi_{OL} \)-critical if \( \chi_{OL}(G - v) < \chi_{OL}(G) \) for all \( v \in V(G) \). The following also hold with \( \chi \) or \( \chi_{L} \) in place of \( \chi_{OL} \).

Corollary 1.4. For any \( \chi_{OL} \)-critical graph \( G \), we have \( \chi_{OL}(G) < \Delta(G) + 2 - \frac{\text{mic}(G)}{|G|} \).

Since we always have \( \text{mic}(G) \geq \frac{2\|G\|}{\chi(G)} \), this gives the following (where \( d(G) \) is the average degree of \( G \)).
Corollary 1.5. For any $\chi_{OL}$-critical graph $G$, we have $\chi_{OL}(G) < \Delta(G) + 2 - \frac{d(G)}{\chi(G)}$.

The regular case of Corollary 1.4 simplifies somewhat.

Corollary 1.6. For any regular $\chi_{OL}$-critical graph $G$, we have $\chi_{OL}(G) < \frac{|G| - \alpha(G)}{|G|} \Delta(G) + 2$.

Note that Corollary 1.6 reduces proving Brooks’ theorem to showing that $\alpha(G) \geq \frac{|G|}{\Delta(G)}$ when $K_{\Delta(G)+1} \not\subseteq G$.

A simple probabilistic argument gives a reasonable bound on $\text{mic}(G)$ for triangle-free graphs and we get the following.

Corollary 4.10. Triangle-free graphs are online $(\Delta + 1 - \frac{1}{4} \log(\Delta))$-choosable.

2 Either many edges or an Alon-Tarsi orientable subgraph

Let $T_k$ be the Gallai trees with maximum degree at most $k - 1$, excepting $K_k$. For a graph $G$, let $W_k(G)$ be the $v \in V(G)$ contained in some $K_k$ in $G$.

2.1 Extending Alon-Tarsi orientations

In [15] Kostochka and Stiebitz gave a method for extending list colorings into Gallai trees. We generalize these ideas in terms of extension of orientations.

Lemma 2.1. Let $G$ be a multigraph without loops and $f: V(G) \to \mathbb{N}$. If there are $F \subseteq G$ and $Y \subseteq V(G)$ such that:

1. any multiple edges in $G$ are contained in $G[Y]$; and
2. $f(v) \geq d_G(v)$ for all $v \in V(G) - Y$; and
3. $f(v) \geq d_G[Y](v) + d_F(v) + 1$ for all $v \in Y$; and
4. For each component $T$ of $G - Y$ there are different $x_1, x_2 \in V(T)$ where $N_T[x_1] = N_T[x_2]$ and $T - \{x_1, x_2\}$ is connected such that either:

   (a) there are $x_1y_1, x_2y_2 \in E(F)$ where $y_1 \neq y_2$ and $N(x_i) \cap Y = \{y_i\}$ for $i \in [2]$; or
   (b) $|N(x_2) \cap Y| = 0$ and there is $x_1y_1 \in E(F)$ where $N(x_1) \cap Y = \{y_1\}$,

then $G$ is $f$-AT.

Proof. Suppose not and choose a counterexample $(G, f, F, Y)$ minimizing $|G - Y|$. If $|G - Y| = 0$, then $f(v) \geq d_G(y) + 1$ for all $v \in V(G)$, so any acyclic orientation of $G$ will do. Hence $|G - Y| > 0$.

Pick a component $T$ of $G - Y$ and pick $x_1, x_2 \in V(T)$ as guaranteed by (4). First, suppose (4a) holds. Put $G' := (G - T) + y_1y_2, F' := F - T, Y' := Y - V(T)$ and let $f'$ be $f$ restricted to $V(G')$. Then $G'$ has an orientation $D'$ where $f'(v) \geq d^{-1}_{D'}(v) + 1$ for all $v \in V(D')$ and $EE(D') \neq EO(D')$, for otherwise $(G', f', F', Y')$ would contradict minimality. By symmetry we may assume that the new edge $y_1y_2$ is directed toward $y_2$. Now we use the orientation
of $D'$ to construct the desired orientation of $D$. First, we use the orientation on $D' - y_1y_2$ on $G - T$. Now, order the vertices of $T$ as $x_1, x_2, z_1, z_2, \ldots$ so that every vertex has at least one neighbor to the right. Orient the edges of $T$ left-to-right in this ordering. Finally, we use $y_1x_1$ and $x_2y_2$ and orient all other edges between $T$ and $G - T$ away from $T$. Plainly, $f(v) \geq d^+_D(v) + 1$ for all $v \in V(D)$. Now, there is a bijection between eulerian subgraphs of $D'$ containing $y_1y_2$ and eulerian subgraphs of $D - (T - \{x_1, x_2\})$ containing $y_1x_1x_2y_2$, moreover this bijection preserves the parity of the eulerian subgraph. Any other eulerian subgraph $A$ of $D$ must contain $y_1x_1$ and not $x_2y_2$. Since $x_1$ must have in-degree 1 in $A$, it must also have out-degree 1 in $A$. We show that $A$ has a mate $A'$ of opposite parity. Suppose $x_2 \not\in A$ and $x_1z_1 \in A$; then we make $A'$ by removing $x_1z_1$ from $A$ and adding $x_2z_1$. If $x_2 \in A$ and $x_1x_2z_1 \in A$, we make $A'$ by removing $x_1x_2z_1$ and adding $x_1z_1$. We conclude that $EE(D) \neq EO(D)$, a contradiction.

Now suppose (4b) holds. Put $G' := G - T$, $F' := F - T$, $Y' := Y - V(T)$ and define $f'$ by $f'(v) = f(v)$ for all $v \in V(G' - y_1)$ and $f'(y_1) = f(y_1) - 1$. Then $G'$ has an orientation $D'$ where $f'(v) \geq d^+_D(v) + 1$ for all $v \in V(D')$ and $EE(D') \neq EO(D')$, for otherwise $(G', f', F', Y')$ would contradict minimality. We orient $G - T$ according to $D$, orient $T$ as in the previous case, again use $y_1x_1$ and orient all other edges between $T$ and $G - T$ away from $T$. Since we decreased $f(y_1)$ by 1, the extra out edge of $y_1$ is accounted for and we have $f(v) \geq d^+_D(v) + 1$ for all $v \in V(D)$. Again any additional eulerian subgraph must contain $y_1x_1$ and since $x_2$ has no neighbor in $G - T$ we can use $x_2$ as before to build a mate of opposite parity for any additional eulerian subgraph. Hence $EE(D) \neq EO(D)$ giving our final contradiction.

Lemma 2.2. Let $r \geq 0$, $k \geq r + 4$ and $G \neq K_k$ be a graph with $x \in V(G)$ such that:

1. $G - x$ is connected; and
2. $d_G(x) \geq r + 2$; and
3. $|N(x) \cap W^k(G - x)| \geq 1$; and
4. $d_G(v) \leq k - 1$ for all $v \in V(G - x)$.

Then $G$ is $f$-AT where $f(x) = d_G(x) - r$ and $f(v) = d_G(v)$ for all $v \in V(G - x)$.

Proof. Suppose not and choose a counterexample minimizing $|G|$. Let $Q$ be the non-separating vertices in $G - x$. Suppose we have $y \in Q$ such that $G - y$ satisfies all the hypotheses of the theorem. Then minimality of $|G|$ shows that $G - y$ is $f'$-AT where $f'(v) := f(v) + d_{G - y}(v) - d_G(v)$ for $v \in V(G)$. Create an orientation $D$ of $G - y$ by directing all edges incident to $y$ into $y$. These new edges are on no cycle and thus the eulerian subgraph counts did not change. Also, we have increased the out degree of any vertex $v$ by at most $d_G(v) - d_{G - y}(v)$. Hence $G$ is $f$-AT, a contradiction. Therefore $G - y$ must fail some hypothesis for each $y \in Q$; note that it is only possible for $G - y$ to fail (2) or (3).

We show that $Q \subseteq N(x)$. Suppose otherwise that we have $y \in Q - N(x)$. Since (2) is satisfied for $G - y$, (3) must fail and hence $y$ is contained in a $K_{k-1}$, call it $B$, in $G - x$ such that $N(x) \cap B \neq \emptyset$. Pick $z \in N(x) \cap B$. Since $d_G(z) \leq k - 1$ we must have
Lemma 2.3 (Rubin [7]). The following lemma from [7] (this is often referred to as “Rubin’s Block Theorem”).

Theorem. Let \(G\) be a graph. Then \(G\) is \(d\)-choosable if and only if for every subset \(A\) of \(V(G)\) there is a \(d\)-list coloring of \(G\) for every subset of \(A\).

Lemma 2.4. For a connected graph \(G\), the following are equivalent:

1. \(G\) is not a Gallai tree,
2. \(G\) contains an induced even cycle with at most one chord,
3. \(G\) is \(d_0\)-choosable,
4. \(G\) is \(d_0\)-\(AT\),
5. \(G\) has an orientation \(D\) where \(d_G(v) \geq d^+_D(v) + 1\) for all \(v \in V(D)\), \(EE(D) \in \{2, 3\}\) and \(EO(D) \in \{0, 1\}\).

Proof. That (1), (2) and (3) are equivalent is the characterization of \(d_0\)-choosable graphs in [3] and [7]. Since (5) implies (4) and (4) implies (3) it will suffice to show that (2) implies (5).

Proof. If \(A, B\) is a separation of \(G\) such that \(G[A]\) and \(G[B]\) are connected \(d_0\)-\(AT\) graphs and \(A \cap B = \{x\}\), then \(G\) is \(f\)-\(AT\) where \(f(v) = d_G(v)\) for all \(v \in V(G) - x\) and \(f(x) = d_G(x) - 1\).
Proof. By Lemma 2.4 we may choose an orientation $D_A$ of $A$ with $d^+(v) < d(v)$ for all $v \in V(D_A)$ and $EE(D_A) \neq EO(D_A)$ and an orientation $D_B$ of $B$ with $d^+(v) < d(v)$ for all $v \in V(D_B)$ and $EE(D_B) \neq EO(D_B)$. Together these give the desired orientation $D$ of $G$ since no cycle has vertices in both $A - x$ and $B - x$ and thus $EE(D) - EO(D) = EE(D_A)EE(D_B) + EO(D_A)EO(D_B) - (EE(D_A)EO(D_B) + EO(D_A)EE(D_B)) = (EE(D_A) - EO(D_A))(EE(D_B) - EO(D_B)) \neq 0$. □

Lemma 2.6 restricts the interaction of a high vertex and a single low component. Similarly to [15] we’ll use the following lemma to restrict a high vertex’s interaction with two low components.

Lemma 2.6. Let $k \geq 4$ and let $G$ be a graph with $x \in V(G)$ such that:

1. $G - x$ has two components $H_1, H_2 \in \mathcal{T}_k$; and
2. $|N(x) \cap V(H_i)| = 2$ for $i \in [2]$; and
3. $|N(x) \cap W^k(H_i)| \in [2]$ for $i \in [2]$; and

Then $G$ is $f$-AT where $f(x) = d_G(x) - 1$ and $f(v) = d_G(v)$ for all $v \in V(G - x)$.

Proof. Using Lemma 2.5 we just need to show that $Q_i := G \{x \cup V(H_i)\}$ is $d_0$-AT for $i \in [2]$; that is show that $Q_i$ is not a Gallai tree. If $Q_i$ is a Gallai tree, then $x$’s two neighbors in $H_i$ must be in the same block in $H_i$ and this block must be a $K_{k-1}$, but this creates a diamond since $k \geq 4$, impossible. □

Combining Lemma 2.2 and Lemma 2.6 gives the following.

Lemma 2.7. Let $k \geq 5$ and let $G$ be a graph with $x \in V(G)$ such that:

1. $K_k \not\subseteq G$; and
2. $G - x$ has components $H_1, H_2, \ldots, H_t \in \mathcal{T}_k$; and
3. $d_G(v) \leq k - 1$ for all $v \in V(G - x)$; and
4. $|N(x) \cap W^k(H_i)| \geq 1$ for $i \in [t]$; and
5. $d_G(x) \geq t + 2$.

Then $G$ is $f$-AT where $f(x) = d_G(x) - 1$ and $f(v) = d_G(v)$ for all $v \in V(G - x)$.

Proof. Since $d_G(x) \geq t + 2$, either $x$ has 3 neighbors in some $H_i$ or $x$ has two neighbors in each of $H_i, H_j$. In either case, let $C_1, \ldots, C_q$ be the other components of $G - x$. For each $i \in [q]$, pick $z_i \in N(x) \cap V(C_i)$. Then order the vertices of $C_i$ with $z_i$ first and orient all the edges in $C_i$ to the left with respect to this ordering. Now orient all edges between $C_i$ and $G - C_i$ into $C_i$. Note that each vertex in $C_i$ has in-degree at least one and no cycle passes through $C_i$. Hence we can complete the orientation using one of Lemma 2.2 or Lemma 2.6 to get our desired orientation $D$ of $G$. □
To deal with more than one high vertex we need to define the following auxiliary bipartite graph. For a graph $G$, $\{X, Y\}$ a partition of $V(G)$ and $k \geq 4$, let $B_k(X, Y)$ be the bipartite graph with one part $Y$ and the other part the components of $G[X]$. Put an edge between $y \in Y$ and a component $T$ of $G[X]$ iff $N(y) \cap W^k(T) \neq \emptyset$. This lemma gives the substantive improvement over [15] on the lower bound on the number of edges in a list critical graph.

**Lemma 2.8.** Let $k \geq 8$ and let $G$ be a graph with $Y \subseteq V(G)$ such that:

1. $K_k \not\subseteq G$; and
2. the components of $G - Y$ are in $T_k$; and
3. $d_G(v) \leq k - 1$ for all $v \in V(G - Y)$; and
4. with $B := B_k(V(G - Y), Y)$ we have $\delta(B) \geq 3$.

Then $G$ is $f$-AT where $f(y) = d_G(y) - 1$ for $y \in Y$ and $f(v) = d_G(v)$ for all $v \in V(G - Y)$.

**Proof.** Suppose not and choose a counterexample $G$ minimizing $|G|$. Suppose a component of $G - Y$ has an endblock $B \neq K_{k-1}$. Let $X_B$ be the non-separating vertices of $B$. Then $G' := G - X_B$ still satisfies the hypotheses of the theorem since the degrees in $B$ are not affected. Hence, by minimality of $|G|$, we have that $G'$ is $f'$-AT where $f'(v) := f(v) + d_G'(v) - d_G(v)$ for all $v \in V(G')$. But we can extend the orientation of $G'$ to $G$ by ordering the vertices of $B$ with the separating vertices before the non-separating vertices, orienting all edges of $B$ to the right and then orienting all edges from $X_B$ to $G - B$ into $X_B$. In doing so, we don’t change the eulerian subgraph counts since no cycle passes through $X_B$. Moreover, we don’t break the out-degree condition since the out-degree of any $v \in V(G)$ as increased by at most $d_G'(v) - d_G(v)$. Therefore $G$ is $f$-AT, a contradiction. Hence every endblock in $G - Y$ is a $K_{k-1}$. We can do the same thing with a $K_{k-1}$ endblock in which the non-separating vertices have no neighbors in $Y$. Thus every endblock $B$ has a non-separating vertex $x_B$ with a neighbor $y_B \in Y$.

To each component $T$ of $G - Y$ we associate a set of edges $u(T)$ from $W^k(T)$ to $Y$. We’ll also classify the components as either type 1, type 2a, type 2b, type 2c or type 3. Call $T$ saturated if it has an endblock $B_T$ such that every non-separating $z \in V(B_T)$ has a neighbor in $Y$.

Suppose $T$ is saturated. If $T$ has one block, we let $u(T)$ be all edges from $T$ to $Y$. Such a $T$ is of type 2a. Otherwise, $T$ has an endblock $B'$ other than $B_T$. First, suppose every non-separating $z \in V(B')$ has a neighbor in $Y$. Then we let $u(T)$ be all edges from non-separating vertices in $B_T$ to $Y$ together with all edges from non-separating vertices in $B'$ to $Y$. Such a $T$ is of type 3. Otherwise some non-separating $v \in V(B')$ has no edge to $Y$ and we let $u(T)$ be all edges from non-separating vertices in $B_T$ to $Y$ together with $v y_B$. Such a $T$ is of type 2b.

Now suppose $T$ is not saturated. If $T$ has at least 3 endblocks $B_1, B_2, B_3$, then put $u(T) := \{x_i y_i | i \in [3]\}$. Such a $T$ is of type 1. If $T$ has only one block, then pick 3 edges from $T$ with different endpoints in $Y$. Such a $T$ is also of type 1. Now suppose $T$ has exactly two endblocks $B_1$ and $B_2$. If, for some $i \in [2]$, there are two edges incident to non-separating vertices in $B_i$ with different endpoints in $Y$, then let $u(T)$ be these two edges together with $x_{B_{3-i}} y_{B_{3-i}}$. Such a $T$ is of type 1.
If we haven’t defined \( u(T) \) yet, then \( T \) must contain a block \( B = K_{k-1} \) that is not an endblock since \( d_B(T) \geq 3 \). Moreover, \( B \) must contain a non-separating \( x_B \) which is adjacent to \( y_B \in Y - \{y_B_1,y_B_2\} \). If some non-separating \( v \in V(B) \) has no edge to \( Y \), then put \( u(T) := \{x_B,y_B_1,x_B,y_B_2,x_B,y_B\} \). In this case, \( T \) is of type 1. Otherwise, let \( u(T) \) be \( \{x_B,y_B_1,x_B,y_B_2\} \) together with all edges from non-separating vertices of \( B \) to \( Y \). Note that as \( T \) has exactly 2 endblocks, \( B \) has \( k - 3 \) non-separating vertices. This final kind of \( T \) is of type 2c.

We need to handle one more issue before proceeding with the application of Lemma 2.1. When \( T \) is type 2a, 2b, 2c or 3, it is possible that \( u(T) \) contains edges of the form \( yx_1 \) and \( yx_2 \). To see why this could be a problem, consider a component \( T \) of type 2a. We’d like to get the degree of \( T \) triples of edges. If \( T \) is type 2b or 2c, then \( T \) has at most \( 6 \) nontricky vertices in any given \( T \), so we don’t need to worry about triples of edges. If \( T \) is of type 2a, 2b, or 2c, we call \( T \) tricky if there are different \( y_1,y_3 \in Y \) and \( x_1,x_2,x_3,x_4 \in V(T) \) such that \( y_1x_1,y_1x_2,y_2x_3,y_2x_4 \in u(T) \). If \( T \) is of type 3, we call \( T \) tricky if there are different \( y_1,y_2 \in Y \) and different \( x_1,x_2,x_3,x_4,x_5,x_6 \in V(T) \) such that \( y_1x_1,y_1x_2,y_2x_3,y_2x_4,y_3x_5,y_3x_6 \in u(T) \).

Now we will create a bipartite graph \( B' \) from \( B \) by splitting vertices corresponding to components \( T \) of \( G - Y \) as follows. Split each nontricky \( T \) into \( T \) and \( T' \) by divvying up \( T \)'s edges in \( B \) so that \( T \) is adjacent to all and only the endpoints in \( Y \) of the edges in \( u(T) \) and \( T' \) is adjacent to the rest. We call \( T \) an original and \( T' \) a copy. We have the nontricky vertex degrees under control since if \( T \) is nontricky type 2a, we have \( d_{B'}(T) \geq k - 1 - 1 = k - 2 \). If \( T \) is nontricky type 2b, then \( d_{B'}(T) \geq k - 2 - 1 + 1 = k - 2 \). If \( T \) is nontricky type 2c, then \( d_{B'}(T) \geq k - 3 + 1 + 1 = k - 2 \). Finally, if \( T \) is nontricky type 3, then \( d_{B'}(T) \geq 2(k - 2) - 2 = 2k - 6 \).

We aim to find an orientation of \( B' \) such that each \( y \in Y \) has in degree at least 2, each original \( T \) of type 1 has in degree at least 1, each original nontricky \( T \) of type 2a, 2b or 2c has in degree at least 2 and each original nontricky \( T \) of type 3 has in degree at least 3. Note that we are not placing any conditions on the copies and the tricky vertices. Let’s see why we want such an orientation. Suppose we have one. We will construct \( F \subseteq G \) needed in Lemma 2.1. For each directed edge \( yT \), pick an edge in \( G \) from \( y \) to \( T \) from \( u(T) \) and add it to \( F \). We claim that these edges satisfy (4) of Lemma 2.1. If \( T \) is type 1 then the selected edge satisfies (4b). If \( T \) is of type 2a, then we have satisfied (4a). If \( T \) is type 2b or 2c, then we satisfy (4a) if the two edges end in the same block of \( T \) and (4b) otherwise. If \( T \) is type 3, then two of the three edges must end in the same block and we satisfy (4a). Also, as each \( y \in Y \) has in degree at least 2, we have \( f(y) = d_G(y) - 1 \geq d_{G_{[Y]}(y)} + d_F(y) + 1 \). It remains to satisfy (4) for the tricky components without breaking (3). Suppose \( T \) is a tricky component of type 2a, 2b, or 2c and \( y_1x_1,y_1x_2,y_2x_3,y_2x_4 \) are in \( u(T) \). If any of these edges are in \( F \) already, remove them. This only makes (3) better and doesn’t change anything else. Now pick one of \( y_1x_1,y_1x_2 \) and one of \( y_2x_3,y_2x_4 \) so that (4) is satisfied (we can by definition of \( u(T) \)) and add them to \( F \). We haven’t broken (3) because \( y_1,y_2 \) had in degree at least two, and \( T \) is still contributing its one in edge. The tricky components of type 3 are handled similarly by picking one edge from each pair. Now \( G \) is \( f \text{-AT} \) by Lemma 2.1 contradiction.

Hence we may assume that \( B' \) has no such orientation. Let \( Q \) be the original vertices of
type 1 in $B'$, $P$ the original nontricky vertices of type 2a, 2b or 2c in $B'$ and $R$ the original nontricky vertices of type 3 in $B'$. Then each $y \in Y$ has degree at least 3 in $B'$, each $T \in Q$ has degree at least 3 in $B'$, each $T \in P$ has degree at least $k - 2$ in $B'$ and each $T \in R$ has degree at least $2k - 6$ in $B'$. Now we apply Lemma 4.1 with $S := Q \cup P \cup R$. Let $H \subseteq B'|S$. For readability, we’ll put $A_H := A \cap V(H)$ for $A \subseteq V(B')$. To get a contradiction, suppose $H$ doesn’t meet the condition in Lemma 4.1. That is, suppose $H$ doesn’t meet the condition in Lemma 4.1. Let $k \in \{3, 4\}$.

Hence $|Q_H| + 2|P_H| + 3|R_H| > |Y_H|$ and $2|Y_H| > 2|Q_H| + (k - 4)|P_H| + (2k - 9)|R_H|$. Hence $2|Q_H| + |P_H| + (2k - 9)|R_H| < 2|Q_H| + 4|P_H| + 6|R_H|$ giving $(k - 8)|P_H| + (2k - 15)|R_H| < 0$ a contradiction since $k \geq 8$.

With a slightly simpler argument we get the following version with asymmetric degree condition on $B$. The point here is that this works for $k \geq 6$. As we’ll see in the next section, the consequence is that we trade a bit in our size bound for the proof to go through with $k \in \{6, 7\}$.

**Lemma 2.9.** Let $k \geq 6$ and let $G$ be a graph with $Y \subseteq V(G)$ such that:

1. $K_k \not\subset G$; and
2. the components of $G - Y$ are in $T_k$; and
3. $d_G(v) \leq k - 1$ for all $v \in V(G - Y)$; and
4. with $B := B_k(V(G - Y), Y)$ we have $d_B(y) \geq 4$ for all $y \in Y$ and $d_B(T) \geq 2$ for all components $T$ of $G - Y$.

Then $G$ is $f$-AT where $f(y) = d_G(y) - 1$ for $y \in Y$ and $f(v) = d_G(v)$ for all $v \in V(G - Y)$.

**Proof.** Suppose not and choose a counterexample $G$ minimizing $|G|$. In the exact same way as the proof of Lemma 2.8 we use minimality of $|G|$ to conclude that if $T$ is a component of $G - Y$, then every endblock of $T$ is $K_{k-1}$ and every endblock $B$ of $T$ has a vertex $x_B$ among its non-separating vertices $X_B$ that has a neighbor $y_B \in Y$.

To each component $T$ of $G - Y$ we associate a set of edges $u(T)$ from $W^k(T)$ to $Y$. We’ll also classify the components as either type 1, type 2a, type 2b or type 3. Call $T$ saturated if it has an endblock $B_T$ such that every non-separating $z \in V(B_T)$ has a neighbor in $Y$.

Suppose $T$ is saturated. If $T$ has one block, we let $u(T)$ be all edges from $T$ to $Y$. Such a $T$ is of type 2a. Otherwise, $T$ has an endblock $B'$ other than $B_T$. First, suppose every non-separating $z \in V(B')$ has a neighbor in $Y$. Then we let $u(T)$ be all edges from nonseparating vertices in $B_T$ to $Y$ together with all edges from non-separating vertices in $B'$ to $Y$. Such a $T$ is of type 3. Otherwise some non-separating $v \in V(B')$ has no edge to $Y$ and we let $u(T)$ be all edges from non-separating vertices in $B_T$ to $Y$ together with $x_{B'}y_{B'}$. Such a $T$ is of type 2b.
Now suppose $T$ is not saturated. If $T$ has at least 2 endblocks $B_1, B_2$, then put $u(T) := \{x_B, y_B \mid i \in [2]\}$. Such a $T$ is of type 1. Otherwise, $T$ has only one block and we pick 2 edges from $T$ with different endpoints in $Y$. Such a $T$ is also of type 1.

We need to handle one more issue before proceeding with the application of Lemma 2.1. When $T$ is type 2a, 2b or 3, it is possible that $u(T)$ contains edges of the form $yx_1$ and $yx_2$. To see why this could be a problem, consider a component $T$ of type 2a. We’d like to get the degree of $T$ in $B$ as large as possible, but it could be that for some $y \in Y$ we have $yx_1$ and $yx_2$ for different $x_1, x_2 \in V(T)$ and hence $T$ has degree at most $k - 2$ in $B$. If this situation happens more than once for $T$, we need to handle $T$ differently. By Lemma 2.7 each $y \in Y$ has at most 2 neighbors in any given $T$, so we don’t need to worry about triples of edges. If $T$ is of type 2a or 2b, we call $T$ tricky if there are different $y_1, y_2 \in Y$ and different $x_1, x_2, x_3, x_4 \in V(T)$ such that $y_1 x_1, y_1 x_2, y_2 x_3, y_2 x_4 \in u(T)$. If $T$ is of type 3, we call $T$ tricky if there are different $y_1, y_2, y_3 \in Y$ and different $x_1, x_2, x_3, x_4, x_5, x_6 \in V(T)$ such that $y_1 x_1, y_1 x_2, y_2 x_3, y_2 x_4, y_3 x_5, y_3 x_6 \in u(T)$.

Now we will create a bipartite graph $B'$ from $B$ by splitting vertices corresponding to components $T$ of $G - Y$ as follows. Split each nontricky $T$ into $T$ and $T'$ by dividing up $T$’s edges in $B$ so that $T$ is adjacent to all and only the endpoints in $Y$ of the edges in $u(T)$ and $T'$ is adjacent to the rest. We call $T$ an original and $T'$ a copy. We have the nontricky vertex degrees under control since if $T$ is nontricky type 2a, we have $d_B(T) \geq k - 1 - 1 = k - 2$. If $T$ is nontricky type 2b, then $d_B(T) \geq k - 2 - 1 + 1 = k - 2$. Finally, if $T$ is nontricky type 3, then $d_B(T) \geq 2(k - 2) - 2 = 2k - 6$.

We aim to find an orientation of $B'$ such that each $y \in Y$ has in degree at least 2, each original $T$ of type 1 has in degree at least 1, each original nontricky $T$ of type 2a or 2b has in degree at least 2 and each original nontricky $T$ of type 3 has in degree at least 3. Note that we are not placing any conditions on the copies and the trickies. Let’s see why we want such an orientation. Suppose we have one. We will construct $F \subseteq G$ needed in Lemma 2.1. For each directed edge $yT$, pick an edge in $G$ from $y$ to $T$ from $u(T)$ and add it to $F$. We claim that these edges satisfy (4) of Lemma 2.1. If $T$ is type 1 then the selected edge satisfies (4b). If $T$ is of type 2a, then we have satisfied (4a). If $T$ is type 2b, then we satisfy (4a) if the two edges end in the same block of $T$ and (4b) otherwise. If $T$ is type 3, then two of the three edges must end in the same block and we satisfy (4a). Also, as each $y \in Y$ has in degree at least 2, we have $f(y) = d_G(y) - 1 \geq d_G[V](y) + d_F(y) + 1$. It remains to satisfy (4) for the tricky components without breaking (3). Suppose $T$ is a tricky component of type 2a or 2b and $y_1 x_1, y_1 x_2, y_2 x_3, y_2 x_4$ are in $u(T)$. If any of these edges are in $F$ already, remove them. This only makes (3) better and doesn’t change anything else. Now pick one of $y_1 x_1, y_1 x_2$ and one of $y_2 x_3, y_2 x_4$ so that (4) is satisfied (we can by definition of $u(T)$) and add them to $F$. We haven’t broken (3) because $y_1, y_2$ had in degree at least two, and $T$ is still contributing its one in edge. The tricky components of type 3 are handled similarly by picking one edge from each pair. Now $G$ is $f$-AT by Lemma 2.1, a contradiction.

Hence we may assume that $B'$ has no such orientation. Let $Q$ be the original vertices of type 1 in $B'$, $P$ the original nontricky vertices of type 2a or 2b in $B'$ and $R$ the original nontricky vertices of type 3 in $B'$. Then each $y \in Y$ has degree at least 3 in $B'$, each $T \in Q$ has degree at least 3 in $B'$, each $T \in P$ has degree at least $k - 2$ in $B'$ and each $T \in R$ has degree at least $2k - 6$ in $B'$. Now we apply Lemma 4.1 with $S := Q \cup P \cup R$. Let $H \triangleq B'[S]$. For readability, we’ll put $A_H := A \cap V(H)$ for $A \subseteq V(B')$. To get a contradiction, suppose
Lemma 2.10. As proved in [15], a computation gives the following.

\[ \|H\|^+ < 2|Y_H| + |Q_H| + 2|P_H| + 3|R_H|. \]

We also have

\[ \|H\|^+ \geq \max \{4|Y_H|, 2|Q_H| + (k - 2)|P_H| + (2k - 6)|R_H|\}. \]

Hence \(|Q_H| + 2|P_H| + 3|R_H| > 2|Y_H|\) and \(2|Y_H| > |Q_H| + (k - 4)|P_H| + (2k - 9)|R_H|\). Hence \(|Q_H| + (k - 4)|P_H| + (2k - 9)|R_H| < |Q_H| + 2|P_H| + 3|R_H|\) giving \((k - 6)|P_H| + 2(k - 6)|R_H| < 0\) a contradiction since \(k \geq 6\).

2.2 Many edges

The rest of the proof is basically taken verbatim from [15]. We need the following definitions:

\[
\begin{align*}
\mathcal{L}_k(G) &:= G[x \in V(G) \mid d_G(x) = k - 1], \\
\mathcal{H}_k(G) &:= G[x \in V(G) \mid d_G(x) \geq k], \\
\sigma_k(G) &:= \left(k - 2 + \frac{2}{k - 1}\right)|\mathcal{L}_k(G)| - 2\|\mathcal{L}_k(G)\|, \\
\tau_{k,c}(G) &:= 2\|\mathcal{H}_k(G)\| + \left(k - c - \frac{2}{k - 1}\right)\sum_{y \in V(\mathcal{H}_k(G))} (d_G(y) - k), \\
\alpha_k &:= \frac{1}{2} - \frac{1}{(k - 1)(k - 2)}, \\
q_k(G) &:= \alpha_k \sum_{v \in V(G) - W_k(G)} (k - 1 - d_G(v)).
\end{align*}
\]

As proved in [15], a computation gives the following.

Lemma 2.10. Let \(G\) be a graph with \(\delta := \delta(G) \geq 3\) and \(0 \leq c \leq \delta + 1 - \frac{2}{\delta}\). If \(\sigma_{\delta + 1}(G) + \tau_{\delta + 1,c}(G) \geq c|\mathcal{H}_{\delta + 1}(G)|\), then \(2\|G\| \geq g_{\delta + 1}(|G|, c)\).

We need the following degeneracy lemma.

Lemma 2.11. Let \(G\) be a graph and \(f : V(G) \rightarrow \mathbb{N}\). If \(\|G\| > \sum_{v \in V(G)} f(v)\), then \(G\) has an induced subgraph \(H\) such that \(d_H(v) > f(v)\) for each \(v \in V(H)\).

Proof. Suppose not and choose a counterexample \(G\) minimizing \(|G|\). Then \(|G| \geq 3\) and we have \(x \in V(G)\) with \(d_G(x) \leq f(x)\). But now \(\|G - x\| > \sum_{v \in V(G - x)} f(v)\), contradicting minimality of \(|G|\).

We’ll also need the following consequence of a lemma in [15] giving a lower bound on \(\sigma_k(T)\) for \(T \in \mathcal{T}_k\).

Lemma 2.12. Let \(k \geq 6\) and \(T \in \mathcal{T}_k\). If \(K_{k-1} \subseteq T\), then \(\sigma_k(T) \geq 2 + q_k(T)\); otherwise \(\sigma_k(T) \geq 2 - \alpha_k + q_k(T)\).
We are now ready to prove the main theorem.

**Theorem 2.13.** Let $G$ be a graph with $\delta := \delta(G) \geq 5$ and $K_{\delta+1} \not\subseteq G$. If $G$ does not have a nonempty induced subgraph $H$ which is $f_{H}-AT$ where $f_{H}(v) := \delta + d_{H}(v) - d_{G}(v)$ for all $v \in V(H)$, then $2\|G\| \geq \kappa_{\delta+1}(|G|, c)$ where $c := (\delta - 2)\alpha_{\delta+1}$ when $\delta \geq 7$ and $c := (\delta - 3)\alpha_{\delta+1}$ when $\delta \in \{5, 6\}$.

**Proof.** Suppose $G$ does not have a nonempty induced subgraph $H$ which is $f_{H}-AT$ where $f_{H}(v) := \delta + d_{H}(v) - d_{G}(v)$ for all $v \in V(H)$. Put $k := \delta + 1$, $L := \mathcal{L}_{k}(G)$ and $\mathcal{H} := \mathcal{H}_{k}(G)$. Using Lemma 2.10 we just need to show that $\sigma_{k}(G) + \tau_{k,c}(G) \geq c |\mathcal{H}|$. Put $W := W^{k}(\mathcal{L})$, $L' := V(\mathcal{L}) - W$ and $H' := \{v \in V(\mathcal{H}) \mid d_{G}(v) = k\}$. For $y \in V(\mathcal{H})$, put $\tau_{k,c}(y) := d_{H}(y) + (k - c + \frac{2}{k-1}) (d_{G}(y) - k)$. We have $\tau_{k,c}(G) = \sum_{y \in V(\mathcal{H})} \tau_{k,c}(y) \geq \sum_{y \in H'} d_{H}(y) + (k - c - \frac{2}{k-1}) |\mathcal{H} - H'| \geq \sum_{y \in H'} (d_{H}(y) + c |\mathcal{H} - H'|$. Thus it will be sufficient to prove that $S := \sigma_{k}(G) + \sum_{y \in H'} d_{H}(y) \geq c |H'|$.

Let $D$ be the components of $\mathcal{L}$ containing $K_{k-1}$ and $C$ the components of $\mathcal{L}$ not containing $K_{k-1}$. Then $D \cup C \subseteq T_{k}$ for otherwise some $T \in D \cup C$ is $d_{0}-AT$ and hence $f_{T}-AT$ and (1) holds. We have $\sigma_{k}(G) = \sum_{T \in D} \sigma_{k}(T) + \sum_{T \in C} \sigma_{k}(T) \geq 2 |\mathcal{D}| + (2 - \alpha_{k}) |\mathcal{C}| + \alpha_{k} \sum_{v \in L'} (k - 1 - d_{\mathcal{L}}(v))$.

Now we define an auxiliary bipartite graph $F$ with parts $A$ and $B$ where:

1. $B = H'$ and $A$ is the disjoint union of the following sets $A_{1}, A_{2}$ and $A_{3}$,

2. $A_{1} = D$ and each $T \in D$ is adjacent to all $y \in H'$ where $N(y) \cap W^{k}(T) \neq \emptyset$,

3. For each $v \in L'$, let $A_{2}(v)$ be a set of $|N(v) \cap H'|$ vertices connected to $N(v) \cap H'$ by a matching in $F$. Let $A_{2}$ be the disjoint union of the $A_{2}(v)$ for $v \in L'$,

4. For each $y \in H'$, let $A_{3}(y)$ be a set of $d_{\mathcal{H}}(y)$ vertices which are all joined to $y$ in $F$. Let $A_{3}$ be the disjoint union of the $A_{3}(y)$ for $y \in H'$.

**Case 1.** $\delta \geq 7$.

Define $f : V(F) \rightarrow \mathbb{N}$ by $f(v) = 1$ for all $v \in A_{2} \cup A_{3}$ and $f(v) = 2$ for all $v \in B \cup A_{1}$. First, suppose $\|F\| > \sum_{v \in V(F)} f(v)$. Then by Lemma 2.11 $F$ has an induced subgraph $Q$ such that $d_{Q}(v) > f(v)$ for each $v \in V(Q)$. In particular, $V(Q) \subseteq B \cup A_{1}$ and $\delta(Q) \geq 3$. Put $Y := B \cap V(Q)$ and let $X$ be $\bigcup_{T \in V(Q) \cap A_{1}} V(T)$. Now $H := G[X \cup Y]$ satisfies the hypotheses of Lemma 2.8 a contradiction.

Therefore $\|F\| \leq \sum_{v \in V(F)} f(v) = 2(|H'| + |\mathcal{D}| - |A_{2}| - |A_{3}|$. By Lemma 2.7, for each $y \in B$ we have $d_{F}(y) \geq k - 1$. Hence $\|F\| \geq (k - 1 - |H'|$. This gives $(k - 3) |H'| \leq 2 |\mathcal{D}| + |A_{2}| + |A_{3}|$. By our above estimate we have $S \geq 2 |\mathcal{D}| + \alpha_{k} \sum_{v \in L'} (k - 1 - d_{\mathcal{L}}(v)) + \sum_{y \in H'} d_{\mathcal{H}}(y) = 2 |\mathcal{D}| + \alpha_{k} |A_{2}| + |A_{3}| \geq \alpha_{k} (2 |\mathcal{D}| + |A_{2}| + |A_{3}|)$. Hence $S \geq \alpha_{k} (k - 3) |H'|$. Thus our desired bound holds by Lemma 2.10.

**Case 2.** $\delta \in \{5, 6\}$.

Define $f : V(F) \rightarrow \mathbb{N}$ by $f(v) = 1$ for all $v \in A_{1} \cup A_{2} \cup A_{3}$ and $f(v) = 3$ for all $v \in B$. First, suppose $\|F\| > \sum_{v \in V(F)} f(v)$. Then by Lemma 2.11 $F$ has an induced subgraph $Q$ such that $d_{Q}(v) > f(v)$ for each $v \in V(Q)$. In particular, $V(Q) \subseteq B \cup A_{1}$ and $d_{Q}(v) \geq 4$ for $v \in B \cap V(Q)$ and $d_{Q}(v) \geq 2$ for $v \in A_{1} \cap V(Q)$. Put $Y := B \cap V(Q)$ and
let $X$ be $\bigcup_{T \in V(T)} V(T)$. Now $H := G[X \cup Y]$ satisfies the hypotheses of Lemma 2.9, a contradiction.

Therefore $\|F\| \leq \sum_{v \in V(F)} f(v) = 3 |H'| + |D| + |A_2| + |A_3|$. By Lemma 2.7 for each $y \in B$ we have $d_F(y) \geq k - 1$. Hence $\|F\| \geq (k - 1) |H'|$. This gives $(k - 4) |H'| \leq |D| + |A_2| + |A_3|$. By our above estimate we have $S \geq 2 |D| + \alpha_k \sum_{v \in L} (k - 1 - d_L(v)) + \sum_{y \in H'} d_H(y) = 2 |D| + \alpha_k |A_2| + |A_3| \geq \alpha_k (|D| + |A_2| + |A_3|)$. Hence $S \geq \alpha_k (k - 4) |H'|$. Thus our desired bound holds by Lemma 2.10.

We note a corollary of the above proof that will be useful in our computations in the section on Ore degrees. When $\mathcal{H}_k(G)$ is edgeless, $A_3$ is empty and $S = \sigma_k(G)$. Also from the proof, we have $\sigma_k(G) \geq 2 |D| + (2 - \alpha_k) |C| + \alpha_k \sum_{v \in L} (k - 1 - d_L(v)) \geq \alpha_k (2 |D| + |A_2|) + 2 (1 - \alpha_k) |C|$. We write $c(G)$ for the number of components of $G$. Since $(2 - \alpha_k) \geq 2 (1 - \alpha_k)$, we have $\sigma_k(G) \geq (k - 3) \alpha_k |\mathcal{H}_k(G)| + 2 (1 - \alpha_k) c(\mathcal{L}(G))$.

**Corollary 2.14.** Let $G$ be a graph with $\delta := \delta(G) \geq 7$ and $K_{\delta+1} \not\subseteq G$ such that $\mathcal{H}_{\delta+1}(G)$ is edgeless. If $G$ does not have a nonempty induced subgraph $H$ which is $f_H$-AT where $f_H(v) := \delta + d_H(v) - d_G(v)$ for all $v \in V(H)$, then $\sigma_{\delta+1}(G) \geq (\delta - 2) \alpha_{\delta+1} |\mathcal{H}_{\delta+1}(G)| + 2 (1 - \alpha_{\delta+1}) c(\mathcal{L}(G))$.

## 3 The bounds for critical graphs

### 3.1 List critical graphs

**Theorem 3.1.** For $k \geq 6$ and $G \neq K_k$ an $L$-critical graph where $L$ is a $(k - 1)$-assignment on $G$, we have $2 \|G\| \geq g_k(|G|, c)$ where $c := (k - 3) \alpha_k$ when $k \geq 8$ and $c := (k - 4) \alpha_k$ when $k \in \{6, 7\}$.

**Proof.** Since $G$ is $L$-critical, we have $\delta(G) \geq k - 1 \geq 5$. If $\delta(G) \geq k$, then $2 \|G\| \geq k \|G\| \geq g_k(|G|, k)$ and we are done. Hence we may assume that $\delta(G) = k - 1$. Since $G \neq K_k$ and $G$ is $L$-critical, we have $K_{\delta(G)+1} \not\subseteq G$.

Suppose $G$ has a nonempty induced subgraph $H$ which is $f_H$-AT where $f_H(v) := \delta + d_H(v) - d_G(v)$ for all $v \in V(H)$. Let $\pi$ be a coloring of $G - H$ from $L$ and let $L'$ be the list assignment on $H$ defined by $L'(v) := L(v) - \pi(N(v) \cap V(G - H))$ for $v \in V(H)$. Then $|L'(v)| \geq |L(v)| - (d_G(v) - d_H(v)) = k - 1 + d_H(v) - d_G(v)$. By Lemma 2.13 $H$ is $f_H$-choosable and hence $H$ is $L'$-colorable. Therefore $G$ is $L$-colorable, a contradiction.

Now applying Lemma 2.13 proves the theorem.

### 3.2 Online list critical graphs

We will need the following lemma from [21] allowing us to patch together online list colorability of parts into online list colorability of the whole.

**Lemma 3.2.** Let $G$ be a graph and $f: V(G) \rightarrow \mathbb{N}$. If $H$ is an induced subgraph of $G$ such that $G - H$ is online $f|_{V(G - H)}$-choosable and $H$ is online $f_H$-choosable where $f_H(v) := f(v) + d_H(v) - d_G(v)$, then $G$ is online $f$-choosable.

Note that applying Lemma 3.2 where $H$ has a single vertex shows that $\delta(G) \geq k - 1$ for an online $f$-critical graph $G$ where $f(v) := k - 1$ for all $v \in V(G)$. 


Theorem 3.3. For \( k \geq 6 \) and \( G \neq K_k \) an online \( f \)-critical graph where \( f(v) := k - 1 \) for all \( v \in V(G) \), we have \( 2 \|G\| \geq g_k(|G|, c) \) where \( c := (k - 3)\alpha_k \) when \( k \geq 8 \) and \( c := (k - 4)\alpha_k \) when \( k \in \{6, 7\} \).

Proof. Since \( G \) is online \( f \)-critical, we have \( \delta(G) \geq k - 1 \geq 5 \). If \( \delta(G) \geq k \), then \( 2 \|G\| \geq k |G| \geq g_k(|G|, k) \) and we are done. Hence we may assume that \( \delta(G) = k - 1 \). Since \( G \neq K_k \) and \( G \) is \( f \)-critical, we have \( K_{k(G)+1} \not\subseteq G \).

Suppose \( G \) has a nonempty induced subgraph \( H \) which is \( f_H \)-AT where \( f_H(v) := \delta(G) + d_H(v) - d_G(v) \) for all \( v \in V(H) \). Let \( \pi \) be a coloring of \( G - H \) from \( L \) and let \( L' \) be the list assignment on \( H \) defined by \( L'(v) := L(v) - \pi(N(v) \cap V(G - H)) \) for \( v \in V(H) \). Then \( |L'(v)| \geq |L(v)| - (d_G(v) - d_H(v)) = k - 1 + d_H(v) - d_G(v) \). By Lemma 1.2, \( H \) is online \( f_H \)-choosable. But then applying Lemma 3.2 shows that \( G \) is online \( f \)-choosable, a contradiction.

Now applying Lemma 2.13 proves the theorem. \( \Box \)

3.3 Alon-Tarsi critical graphs

We need a lemma that serves the same purpose as Lemma 3.2 for orientations.

Lemma 3.4. Let \( G \) be a graph and \( f : V(G) \to \mathbb{N} \). If \( H \) is an induced subgraph of \( G \) such that \( G - H \) is \( f|_{V(G-H)} \)-AT and \( H \) is \( f_H \)-AT where \( f_H(v) := f(v) + d_H(v) - d_G(v) \), then \( G \) is \( f \)-AT.

Proof. Take an orientation of \( G - H \) demonstrating that it is \( f|_{V(G-H)} \)-AT and an orientation of \( H \) demonstrating that it is \( f_H \)-AT. Now orient all the edges between \( H \) and \( G - H \) into \( G - H \). Call the resulting oriented graph \( D \). Then \( D \) satisfies the out degree requirements of being \( f \)-AT since the out degree of the vertices in \( G - H \) haven’t changed and the out degree of each \( v \in V(H) \) has increased by \( d_G(v) - d_H(v) \). Since no directed cycle in \( D \) has vertices in both \( H \) and \( D - H \), the eulerian subgraphs of \( D \) are just all pairings of eulerian subgraphs of \( H \) and \( D - H \). Therefore \( EE(D) - EO(D) = EE(H)EE(D - H) + EO(H)EO(D - H) - (EE(H)EO(D - H) + EO(H)EE(D - H)) = (EE(H) - EO(H))(EE(D - H) - EO(D - H)) \neq 0 \). Hence \( G \) is \( f \)-AT.

Note that applying Lemma 3.4 where \( H \) has a single vertex shows that \( \delta(G) \geq k - 1 \) for a \( k \)-AT-critical graph \( G \).

Theorem 3.5. For \( k \geq 6 \) and \( G \neq K_k \) a \( k \)-AT-critical graph, we have \( 2 \|G\| \geq g_k(|G|, c) \) where \( c := (k - 3)\alpha_k \) when \( k \geq 8 \) and \( c := (k - 4)\alpha_k \) when \( k \in \{6, 7\} \).

Proof. Since \( G \) is \( k \)-AT-critical, we have \( \delta(G) \geq k - 1 \geq 5 \). If \( \delta(G) \geq k \), then \( 2 \|G\| \geq k |G| \geq g_k(|G|, k) \) and we are done. Hence we may assume that \( \delta(G) = k - 1 \). Since \( G \neq K_k \) and \( G \) is \( k \)-AT-critical, we have \( K_{k(G)+1} \not\subseteq G \).

Suppose \( G \) has a nonempty induced subgraph \( H \) which is \( f_H \)-AT where \( f_H(v) := \delta(G) + d_H(v) - d_G(v) \) for all \( v \in V(H) \). Since \( G - H \) is \( (k - 1) \)-AT, applying Lemma 3.4 shows that \( G \) is \( (k - 1) \)-AT, a contradiction.

Now applying Lemma 2.13 proves the theorem. \( \Box \)
4 Good orientations from big cuts

Here we strengthen the lemmas on list colorings and orientations of Kostochka and Yancey [12]. The proofs are identical except we replace their specific method of getting an orientation by the following general lemma. This lemma can be derived from Hall’s theorem via a similar vertex splitting method to that used by Kostochka and Yancey. It also follows by taking an arbitrary orientation and repeatedly reversing paths if doing so gets a gain (really, this is just the proof of the max-flow min-cut theorem). We give the proof for completeness. For \( H \leq G \) put \( \| H \|^{+} := \| H \| + \| H, G - H \| \).

**Lemma 4.1.** Let \( G \) be a graph, \( S \subseteq V(G) \) and \( g: S \rightarrow \mathbb{N} \). Then \( G \) has an orientation such that \( d^{-}(v) \geq g(v) \) for all \( v \in S \) iff for every \( H \leq G[S] \) we have

\[
\| H \|^{+} \geq \sum_{v \in V(H)} g(v).
\]

**Proof.** First, suppose \( G \) has such an orientation and let \( H \leq G[S] \). Then \( \| H \|^{+} = \sum_{v \in V(H)} d^{-}(v) + \| H, G - H \| \geq \sum_{v \in V(H)} d^{-}(v) \geq \sum_{v \in V(H)} g(v) \).

For the other direction, pick an orientation of \( G \) minimizing

\[
\Theta := \sum_{v \in S} \max\{0, g(v) - d^{-}(v)\}.
\]

Note that the orientation has the desired property iff \( \Theta = 0 \). Suppose \( \Theta > 0 \). Then we may choose \( x \in S \) with \( d^{-}(x) < g(x) \). Put

\[
X := \{ v \in V(G) \mid \text{there exists an } xv\text{-directed-path} \}.
\]

For any \( v \in X \) we must have \( v \in S \) and \( d^{-}(v) \leq g(v) \) for otherwise reversing all the edges on an \( xv\)-path violates minimality of \( \Theta \). By definition, all edges between \( X \) and \( G - X \) are directed into \( X \) and hence with \( H := G[X] \) we have

\[
\| H \|^{+} = \sum_{v \in V(H)} d^{-}(v) = d^{-}(x) + \sum_{v \in V(H) - x} d^{-}(v) < \sum_{v \in V(H)} g(v).
\]

This contradicts our supposition and thus completes the proof. \( \square \)

A *kernel* in a digraph \( D \) is an independent set \( I \subseteq V(D) \) such that each vertex in \( V(D) - I \) has an edge into \( I \). A digraph in which every induced subdigraph has a kernel is called *kernel-perfect*. Schauz [21] shows that the well-known Kernel Lemma extends to online choosability.

**Kernel Lemma.** If \( D \) is a kernel-perfect digraph and \( f: V(D) \rightarrow \mathbb{N} \) such that \( f(v) \geq d^{+}(v) + 1 \) for each \( v \in V(D) \), then \( G \) is online \( f \)-choosable.

All bipartite graphs are kernel-perfect, the following lemma from [12] generalizes this fact.
Lemma 4.2. Let $B$ be an independent set in a graph $G$ and put $A := V(G) - B$. Any digraph created from $G$ by replacing each edge in $G[A]$ by a pair of opposite arcs and orienting the edges between $A$ and $B$ arbitrarily is kernel-perfect.

Now we have everything needed to prove the strengthening. For disjoint subsets $A$ and $B$ of $V(G)$, we write $G(A, B)$ for the bipartite subgraph $G[A] \cup B = E(G[A]) - E(G[B])$.

Lemma 4.3. Let $G$ be a graph and $f : V(G) \to \mathbb{N}$ with $f(v) \leq d_G(v) + 1$ for all $v \in V(G)$. If there is independent $A \subseteq V(G)$ such that for each $Q \subseteq G(A, V(G - A))$ we have

$$\|Q\|^+ \geq \sum_{v \in V(Q)} d_G(v) + 1 - f(v).$$

then $G$ is online $f$-choosable.

Proof. Applying Lemma 4.1 on $G(A, V(G - A))$ with $S = V(G)$ and $g(v) := d_G(v) + 1 - f(v)$ for all $v \in V(G)$ gives an orientation of $G$ where $d^-(v) \geq d_G(v) + 1 - f(v)$ for each $v \in V(G)$ and thus $f(v) \geq d^+(v) + 1$. Using this orientation in Lemma 4.2 for the edges between $A$ and $V(G - A)$ then gives the desired kernel-perfect orientation of $G$ showing that $G$ is online $f$-choosable.

Lemma 4.4. Let $G$ be a nonempty graph and $f : V(G) \to \mathbb{N}$ with $f(v) \leq d_G(v) + 1$ for all $v \in V(G)$. If there is independent $A \subseteq V(G)$ such that

$$\|A, G - A\| \geq \sum_{v \in V(G)} d_G(v) + 1 - f(v),$$

then $G$ has a nonempty induced subgraph $H$ that is online $f_H$-choosable where $f_H(v) := f(v) + d_H(v) - d_G(v)$ for $v \in V(H)$.

Proof. Let $A \subseteq V(G)$ be an independent set with $\|A, G - A\| \geq \sum_{v \in V(G)} (d_G(v) + 1 - f(v))$ and put $B := V(G - A)$. Choose nonempty $H \subseteq G$ with $\|A \cap V(H), B \cap V(H)\| \geq \sum_{v \in V(H)} (d_H(v) + 1 - f_H(v))$ minimizing $|H|$. Suppose $H$ is not online $f_H$-choosable. Then, by Lemma 4.3 we have $Q \leq G(A \cap V(H), B \cap V(H))$ with $\|Q\|^+ < \sum_{v \in V(Q)} (d_H(v) + 1 - f_H(v))$. Now $Q \neq H$ by our assumption on $H$, hence $Z := H - Q$ is a nonempty induced subgraph of $G$ with $\|A \cap V(Z), B \cap V(Z)\| = \|A \cap V(H), B \cap V(H)\| - \|Q\|^+ > \sum_{v \in V(H)} (d_H(v) + 1 - f_H(v)) - \sum_{v \in V(Q)} (d_H(v) + 1 - f_H(v)) = \sum_{v \in V(Z)} (d_Z(v) + 1 - f_Z(v))$ contradicting the minimality of $|H|$.

As a special case we get the following lower bound on the number of edges. Recall that $\text{mic}(G)$ is the maximum of $\sum_{v \in I} d_G(v)$ over all independent sets $I$ of $G$.

Theorem 4.5. For any graph $G$ we have either:

1. $G$ has a nonempty induced subgraph $H$ which is online $f_H$-choosable where $f_H(v) := \delta(G) + d_H(v) - d_G(v)$ for all $v \in V(H)$; or

2. $2\|G\| \geq (\delta(G) - 1)|G| + \text{mic}(G) + 1$.

We now give a few applications of these ideas.
4.1 Classification of (online) degree-choosable graphs

Using Theorem 4.3 it is easy to prove Brooks’ theorem for online list coloring given Brooks’ theorem for ordinary coloring; to wit: a minimum counterexample $G$ is $\Delta(G)$-regular, applying Brooks’ theorem gives $\alpha(G) \geq \frac{|G|}{\Delta(G)}$ and hence $\text{mic}(G) \geq |G|$ contradicting Theorem 4.5.

In fact, from Lemma 4.4 we can derive the classification of (online) $d_0$-choosable graphs. The following generalizes this classification.

Theorem 4.6. For a connected graph $G$, we have $\text{mic}(G) \geq |G| - 1$ with equality only when $G$ is a Gallai tree.

Proof. Suppose not and choose a counterexample $G$ minimizing $|G|$. Plainly, we must have $|G| \geq 2$. First, suppose $\text{mic}(G) < |G| - 1$. Let $v$ be a noncutvertex in $G$. By minimality of $|G|$, we have $\text{mic}(G - v) \geq |G| - 2$. Let $A$ be an independent set in $G - v$ with $|A, G - v - A| = \text{mic}(G - v)$. Put $A' := A \cup \{v\}$ if $|v, A| = 0$ and $A' := A$ otherwise. Then $A'$ is independent in $G$ and $|A', G - A'| \geq |G| - 1$, a contradiction.

Now suppose $G$ is not a Gallai tree and $\text{mic}(G) = |G| - 1$. As above we know that for any noncutvertex $v$ we have $\text{mic}(G) \geq \text{mic}(G - v) + 1$; it follows, by minimality of $G$, that $G - v$ is a Gallai tree. So, if $G$ had more than one block, we could remove a vertex from each of two endblocks to show that every block of $G$ is either complete or an odd cycle contradicting our assumption. Hence $G$ is 2-connected and in particular $\delta(G) \geq 2$. Pick $v \in V(G)$. Suppose $G - v$ is 2-connected. Then $G - v$ is complete or an odd cycle. If $G - v$ is complete then taking $v$ together with a nonneighbor of $v$ shows $\text{mic}(G) \geq |G|$. Suppose $G - v$ is an odd cycle. If $d(v) = 2$, then pick a maximum independent set avoiding $N(v)$ together with $v$ to show $\text{mic}(G) \geq |G|$. Otherwise we may pick maximum independent set containing 2 elements of $N(v)$ again showing $\text{mic}(G) \geq |G|$. Hence $G - v$ has a cutvertex for every $v \in V(G)$. Therefore $v$ is adjacent to every noncutvertex in every endblock of $G - v$ and hence $G - v$ must have exactly two endblocks. Since internal noncutvertices could be removed leaving a 2-connected graph, we conclude that $G - v$ must be a path. Since this was for any $v$, $G$ must be an even cycle. But then $\text{mic}(G) \geq |G|$, a contradiction.

It remains to show that $\text{mic}(G) = |G| - 1$ when $G$ is a Gallai tree. This can be proved directly, but we won’t do so since it follows immediately from the construction of bad $d_0$-assignments on Gallai trees and applying Lemma 4.4 as in the proof of the following Corollary.

\[\square\]

Corollary 4.7. A connected graph is online $d_0$-choosable if it is not a Gallai tree.

Proof. Let $G$ be a connected graph that is not a Gallai tree. By Theorem 4.6 we have $\text{mic}(G) \geq |G|$. Applying Lemma 4.4 with $f(v) := d_G(v)$ for all $v \in V(G)$ gives a nonempty induced subgraph $H$ of $G$ that is online $f_H$-choosable where $f_H(v) := f(v) + d_H(v) - d_G(v) = d_H(v)$ for $v \in V(H)$. Hence $G$ is online $d_0$-choosable by Lemma 3.2. \[\square\]

4.2 Online choosability of triangle-free graphs

We write $\log_2(x)$ for the base 2 logarithm of $x$. We can get a reasonably good lower bound on $\text{mic}(G)$ for triangle-free graphs using a simple probabilistic technique of Shearer and its modification by Alon (see [2]).
Lemma 4.8. If $G$ is a triangle-free graph, then $\text{mic}(G) \geq \frac{1}{4} \sum_{v \in V(G)} \log(d(v))$.

Proof. Let $W$ be a random independent set in $G$ chosen uniformly from all independent sets in $G$. For each $v \in V(G)$ put $X_v := d(v) |\{v\} \cap W| + |N(v) \cap W|$. We claim that $E(X_v) \geq \frac{1}{2} \log(d(v))$. This implies the lemma since by linearity of expectation $2 \text{mic}(G) \geq E\left(\sum_{v \in V(G)} X_v\right) \geq \frac{1}{2} \sum_{v \in V(G)} \log(d(v))$.

To prove the claim, let $H$ be the subgraph of $G$ induced on $V(G) - (N(v) \cup \{v\})$, fix and independent set $S$ in $H$ and let $X$ be the set of all nonneighbors of $S$ in $N(v)$. Put $x := |X|$. It will suffice to bound the conditional expectation for each possible $S$ as follows:

$$E(X_v \mid W \cap V(H) = S) \geq \frac{\log(d(v))}{2}.$$ 

For each $S$, there are exactly $2^x + 1$ possibilities for $W$ and we see that the conditional expectation is exactly $\frac{d(v) + 2x - 1}{2^{x+1}}$. Suppose this is less than $\frac{\text{lg}(d(v))}{2}$ for some $x$. Then $2^x \left(\frac{\log(d(v))}{2} - \frac{x}{2}\right) > d(v) - \frac{\text{lg}(d(v))}{2}$. Put $t := \log(d(v)) - x$. We have $\frac{d(v)}{2^x} = \frac{d(v)}{2^t} \left(\frac{\log(d(v))}{2} - \frac{\text{lg}(d(v))}{2} - t\right) > d(v) - \frac{\log(d(v))}{2}$ and hence $\frac{t}{2^t} > 1$, a contradiction. \qed

Theorem 4.9. Let $G$ be a triangle-free graph and define $f : V(G) \to \mathbb{N}$ by $f(v) := d_G(v) + 1 - \left\lfloor \frac{1}{4} \log(d_G(v)) \right\rfloor$. Then $G$ has a nonempty induced subgraph $H$ that is online $f_H$-choosable where $f_H(v) := f(v) + d_H(v) - d_G(v)$ for $v \in V(H)$.

Proof. Immediate upon applying Lemma 4.4 to $G$ since

$$\sum_{v \in V(G)} d_G(v) + 1 - f(v) = \sum_{v \in V(G)} \left\lfloor \frac{1}{4} \log(d_G(v)) \right\rfloor \leq \text{mic}(G).$$ \hspace{1cm} \qed

Corollary 4.10. If $G$ is a triangle-free graph with $\Delta(G) \leq t$ for some $t \in \mathbb{N}$, then $G$ is online $(t + 1 - \left\lfloor \frac{1}{4} \log(t) \right\rfloor)$-choosable.

Proof. Suppose not and choose a counterexample $G$ and $t \in \mathbb{N}$ so as to minimize $|G|$. Put $f(v) := d_G(v) + 1 - \left\lfloor \frac{1}{4} \log(d_G(v)) \right\rfloor$. By Theorem 4.9 $G$ has a nonempty induced subgraph $H$ that is online $f_H$-choosable where $f_H(v) := f(v) + d_H(v) - d_G(v)$ for $v \in V(H)$. Since $t + 1 - \left\lfloor \frac{1}{4} \log(t) \right\rfloor \geq d_G(v) + 1 - \left\lfloor \frac{1}{4} \log(d_G(v)) \right\rfloor$ for all $v \in V(G)$, we have that $H$ is $g(v)$-choosable where $g(v) := t + 1 - \left\lfloor \frac{1}{4} \log(t) \right\rfloor + d_H(v) - d_G(v)$. Now applying minimality of $|G|$ and Lemma 3.2 gives a contradiction. \hspace{1cm} \qed

The best known bounds for the chromatic number of triangle-free graphs are Kostochka’s upper bound of $\frac{2}{3} \Delta + 2$ in [13] (see [19] for a proof in English) for small $\Delta$ and Johansson’s upper bound of $\frac{9\Delta}{\ln(\Delta)}$ for large $\Delta$. Johansson’s proof also works for list coloring, but not for online list coloring. To the best of our knowledge Corollary 4.10 is the best known upper bound for online list colorings of triangle-free graphs. Additionally, Corollary 4.10 improves on Johansson’s bound for list coloring for $\Delta \leq 8000$. The bound can surely be improved by a more complicated computation of $\text{mic}(G)$, but not beyond around $\Delta + 1 - \left\lfloor 2 \log(\Delta) \right\rfloor$ via this method as can be seen by examples of triangle-free graphs with independence number near $\frac{2 \log(\Delta)}{\Delta}$. 

20
4.3 Ore Brooks for maximum degree four

Kostochka and Yancey’s bound [12] shows that if $G$ is 4-critical, then $\|G\| \geq \left\lceil \frac{5\|G\| - 2}{3} \right\rceil$. If we try to analyze 4-critical graphs with edgeless high vertex subgraphs by putting this lower bound on the number of edges together with the results on orientations and list coloring obtained in [12], the bounds miss each other. Using the improved bound from Lemma 5.1 we get an exact bound on the number of edges in such a graph.

**Lemma 4.11.** For a critical graph $G$ with $\Delta(G) \leq \chi(G) = 4$ such that $\mathcal{H}(G)$ is edgeless we have $\|G\| = \left\lceil \frac{5\|G\| - 2}{3} \right\rceil$ and $|G|$ is not a multiple of 3.

**Proof.** Since $G$ is 4-critical, applying Lemma 5.1 gives $2\|G\| < \left(3 + \frac{1}{3}\right)|G| = \frac{10}{3}|G|$. By Kostochka and Yancey’s bound we have $\left\lceil \frac{5\|G\| - 2}{3} \right\rceil \leq \|G\| < \frac{5}{3}|G|$. Hence $\|G\| = \left\lceil \frac{5\|G\| - 2}{3} \right\rceil$ and $|G|$ is not a multiple of 3.

It is easy to see that contracting a diamond in a critical graph $G$ with $\Delta(G) \leq \chi(G) = 4$ such that $\mathcal{H}(G)$ is edgeless gives another such graph. The following characterization of these graphs is natural. We have recently learned that Postle has proven this using an extension of the potential method of Kostochka and Yancey.

**Conjecture 4.12.** Every critical graph $G$ with $\Delta(G) \leq \chi(G) = 4$ such that $\mathcal{H}(G)$ is edgeless, except $K_4$, has an induced diamond. In particular, any such $G$ can be reduced to $K_4$ by a sequence of diamond contractions.

5 Ore Brooks for online list coloring

For a graph $G$, let $\mathcal{H}(G)$ be the subgraph of $G$ induced on the vertices of degree greater than $\delta(G)$ and $L(G)$ the subgraph of $G$ induced on the vertices of degree $\delta(G)$.

**Lemma 5.1.** If $G$ is a graph such that $\mathcal{H}(G)$ is edgeless and $\Delta(G) = \delta(G) + 1$, then either:

1. $G$ has a nonempty induced subgraph $H$ which is online $f_H$-choosable where $f_H(v) := \delta(G) + d_H(v) - d_G(v)$ for all $v \in V(H)$; or
2. $2\|G\| < \left(\delta(G) + \frac{1}{\delta(G)}\right)|G|$.

**Proof.** Put $\delta := \delta(G)$ and suppose $2\|G\| \geq \left(\delta + \frac{1}{\delta}\right)|G|$. Then $|\mathcal{H}(G)| + \delta|G| = 2\|G\| \geq \left(\delta + \frac{1}{\delta}\right)|G|$ and hence $|G| \leq \delta|\mathcal{H}(G)|$. Therefore $\|L(G), \mathcal{H}(G)\| \geq (\delta + 1)|\mathcal{H}(G)| \geq |\mathcal{H}(G)| + |G|$. Plugging into (2) of Lemma 4.5 gives $2\|G\| \geq 2\|G\| + 1$, which is impossible; hence we get the desired induced subgraph $H$ and (1) holds. □

To break up our computations we reformulate Lemma 5.1 as an upper bound on $\sigma$.

**Lemma 5.2.** If $G$ is a graph such that $\mathcal{H}(G)$ is edgeless and $\Delta(G) = \delta(G) + 1$, then either:

1. $G$ has a nonempty induced subgraph $H$ which is online $f_H$-choosable where $f_H(v) := \delta(G) + d_H(v) - d_G(v)$ for all $v \in V(H)$; or
2. $\sigma_{\delta(G)+1}(G) < \left(4 - \frac{2}{\delta(G)}\right)|\mathcal{H}(G)|$.

**Proof.** Put $\delta := \delta(G)$. Suppose (1) does not hold. By Lemma 5.1 we have $2\|G\| < (\delta + \frac{1}{\delta})|G|$. We have $\delta|G| + |\mathcal{H}(G)| = 2\|G\| < (\delta + \frac{1}{\delta})|G|$ giving $|\mathcal{H}(G)| < \frac{|G|}{\delta}$ and hence $|\mathcal{L}(G)| > (\delta - 1)|\mathcal{H}(G)|$. The lemma follows from the following computation:

$$\sigma_{\delta+1}(G) = 2\delta |\mathcal{L}(G)| - 2 |\mathcal{L}(G)| - \left(\delta + 1 - \frac{2}{\delta}\right) |\mathcal{L}(G)|$$

$$= 2 \|G\| - \left(\delta + 1 - \frac{2}{\delta}\right) |\mathcal{L}(G)|$$

$$< \left(\delta + \frac{1}{\delta}\right) |G| - \left(\delta + 1 - \frac{2}{\delta}\right) |\mathcal{L}(G)|$$

$$= -\left(1 - \frac{3}{\delta}\right) |\mathcal{L}(G)| + \left(\delta + \frac{1}{\delta}\right) |\mathcal{H}(G)|$$

$$< -\left(1 - \frac{3}{\delta}\right) (\delta - 1) |\mathcal{H}(G)| + \left(\delta + \frac{1}{\delta}\right) |\mathcal{H}(G)|$$

$$= \left(4 - \frac{2}{\delta}\right) |\mathcal{H}(G)|.$$ 

□

By combining Lemma 5.2 with Corollary 2.14 we can prove the Ore version of Brooks’ theorem for online list coloring for $\Delta \geq 11$. With a bit more work we will improve this to $\Delta \geq 10$. First, we can actually get a bit more out of Lemma 4.5 by considering independent sets of low vertices that have no high neighbors. Such sets can be added to $V(\mathcal{H}(G))$ to get a cut with more edges. To apply this idea we need the following counting lemma. For a graph $G$ and $t \in \mathbb{N}$, let $\beta_t(G)$ be the size of a largest independent set of degree $t$ vertices in $G$; that is, $\beta_t(G) := \alpha(G[x|d(x) = t])$.

**Lemma 5.3.** Fix $k \geq 6$. Let $G$ be a Gallai forest with maximum degree at most $k - 1$. We have the following inequality:

$$(k - 1)\beta_{k-1}(G) + \sum_{v \in V(G)} k - 1 - d(v) \geq \frac{2(k - 3)}{k - 2} |G| - \frac{(k - 1)(k - 4)}{k - 2} c(G).$$

**Proof.** It will suffice to prove that for any Gallai tree with maximum degree at most $k - 1$ we have:

$$(k - 1)\beta_{k-1}(T) + \sum_{v \in V(T)} k - 1 - d(v) \geq \frac{2(k - 3)}{k - 2} |T| - \frac{(k - 1)(k - 4)}{k - 2}.$$

Suppose not and choose a counterexample $T$ minimizing $|T|$. First, if $T$ has only one block it is easy to see that the inequality is satisfied. Let $B$ be an endblock of $T$ and say $x$ is the cutvertex in $B$. Suppose $\chi(B) \leq k - 3$. Put $T' := T - (B - x)$. By minimality of $|T|$, $T'$ satisfies the inequality. When we add $B - x$ back in, the left side increases by
\((k-\chi(B))(|B|-1)-(|B|-1) \geq 2(|B|-1)\). But the right side increases by only \(\frac{2(k-3)}{k-2}(|B|-1)\)
and hence \(T\) is not a counterexample, a contradiction.

Hence \(B\) is either \(K_{k-2}\) or \(K_{k-1}\). Consider \(T' := T-B\). Suppose \(d_{T}(x) = k-1\). Note that none of \(x\)'s neighbors in \(T'\) have degree \(k-1\) in \(T'\) and thus are in no maximum independent set of degree \(k-1\) vertices in \(T'\). Therefore, we can add \(x\) to any such independent set, giving \(\beta_{k-1}(T) > \beta_{k-1}(T')\). Hence, after applying minimality to \(T'\), we see that adding back \(B\) increases the left side by \(k-1 + (k-2) - 1\) if \(B\) is \(K_{k-1}\) and by \(k-1 + 2(k-3) - 2\) if \(B\) is \(K_{k-2}\). Since the right side increases by only \(\frac{2(k-3)}{k-2} |B|\) in both cases, \(T\) satisfies the inequality, a contradiction.

Therefore, it must be that \(B\) is \(K_{k-2}\) and \(d_{T}(x) = k-2\). Now when we add \(B\) back, the left side increases by \(2(k-3) + 1 - 1\) and the right side increases by only \(2(k-3)\) and again \(T\) satisfies the inequality, a contradiction. \(\blacklozenge\)

**Lemma 5.4.** If \(G\) is a graph such that \(\mathcal{H}(G)\) is edgeless and \(\Delta(G) = \delta(G) + 1 \geq 7\), then either:

1. \(G\) has a nonempty induced subgraph \(H\) which is online \(f_{H}\)-choosable where \(f_{H}(v) := \delta(G) + d_{H}(v) - d_{G}(v)\) for all \(v \in V(H)\); or

2. \(|\mathcal{H}(G)| < \frac{\delta(G)(\delta(G)-3)}{\delta(G)-1}c(\mathcal{L}(G))\).

**Proof.** Put \(\delta := \delta(G)\) and suppose (1) does not hold. By Lemma 4.3 we have \(\text{mic}(G) < |\mathcal{L}(G)| + 2|\mathcal{H}(G)| < |\mathcal{L}(G)| + \frac{2}{\delta-1}|\mathcal{L}(G)| = \frac{\delta+1}{\delta-1}|\mathcal{L}(G)|\). But applying Lemma 5.3 to \(\mathcal{L}(G)\) gives \(\text{mic}(G) \geq \frac{2\delta-2}{\delta-1}|\mathcal{L}(G)| - \frac{\delta(\delta-3)}{\delta-1}c(\mathcal{L}(G))\). Combining these inequalities and \(|\mathcal{L}(G)| > (\delta-1)|\mathcal{H}(G)|\) proves the lemma. \(\square\)

**Lemma 5.5.** Every graph \(G\) with \(\delta(G) + 1 = \Delta(G) \geq 10\) such that \(\mathcal{H}(G)\) is edgeless either contains \(K_{\Delta(G)}\) or has a nonempty induced subgraph \(H\) that is online \(f_{H}\)-choosable where \(f_{H}(v) := \Delta(G) - 1 + d_{H}(v) - d_{G}(v)\) for all \(v \in V(H)\).

**Proof.** Suppose not and let \(G\) be a counterexample. Put \(\delta := \delta(G)\). By Corollary 2.14 we have \(\sigma_{\delta+1}(G) \geq (\delta - 2)\alpha_{\delta+1}(G) + 2(1 - \alpha_{\delta+1})c(\mathcal{L}(G))\). By Lemma 5.4 we have \(c(\mathcal{L}(G)) > \frac{(\delta-1)(\delta-5)}{\delta(\delta-3)}|\mathcal{H}(G)|\) and hence. Also, by Lemma 5.2 we have \(\sigma_{\delta+1}(G) < (4 - \frac{2}{\delta})|\mathcal{H}(G)|\). Putting these together, we get \(4 - \frac{2}{\delta} > \alpha_{\delta+1}(\delta - 2) + 2(1 - \alpha_{\delta+1})(\frac{(\delta-1)(\delta-5)}{\delta(\delta-3)})\). But then \(\delta \leq 8\), a contradiction. \(\square\)

**Theorem 5.6.** If \(G\) is a graph with \(\Delta(G) \geq 10\) not containing \(K_{\Delta(G)}\) such that \(\mathcal{H}(G)\) is edgeless, then \(G\) is online \((\Delta(G) - 1)\)-choosable.

**Proof.** Suppose not and choose a counterexample \(G\) minimizing \(|G|\). Then \(G\) is online \(f\)-critical where \(f(v) := \Delta(G) - 1\) for all \(v \in V(G)\). Hence \(\delta(G) \geq \Delta(G) - 1\) and we may apply Lemma 5.5 to get a nonempty induced subgraph \(H\) of \(G\) that is online \(f_{H}\)-choosable where \(f_{H}(v) := \Delta(G) - 1 + d_{H}(v) - d_{G}(v)\) for all \(v \in V(H)\). But then applying Lemma 3.2 shows that \(G\) is \((\Delta(G) - 1)\)-choosable, a contradiction. \(\square\)

Combining Lemma 5.6 with the following version of Brooks’ theorem for online list coloring (first proved in [9]) we get Theorem 5.8.
Lemma 5.7. Every graph with $\Delta \geq 3$ not containing $K_{\Delta+1}$ is online $\Delta$-choosable.

Theorem 5.8. Every graph with $\theta \geq 18$ and $\omega \leq \frac{\theta}{2}$ is online $\left\lfloor \frac{\theta}{2}\right\rfloor$-choosable.

Proof. Suppose not and choose a counterexample $G$ minimizing $|G|$. Put $k := \left\lfloor \frac{\theta(G)}{2}\right\rfloor$. Then $G$ is online $f$-critical where $f(v) := k$ for all $v \in V(G)$. Hence $\delta(G) \geq k$ and thus $\Delta(G) \leq k+1$. If $\Delta(G) = k$, then the theorem follows from Lemma 5.7. Hence we must have $\Delta(G) = k+1$. Therefore $\mathcal{H}(G)$ is edgeless, $\Delta(G) \geq 10$ and $\omega(G) \leq \Delta(G) - 1$. Applying Theorem 5.6 shows that $G$ is online $(\Delta(G) - 1)$-choosable, a contradiction. \qed

The same result for list coloring is an immediate consequence.

Theorem 5.9. Every graph with $\theta \geq 18$ and $\omega \leq \frac{\theta}{2}$ is $\left\lfloor \frac{\theta}{2}\right\rfloor$-choosable.

References

[1] N. Alon and M. Tarsi, Colorings and orientations of graphs, Combinatorica 12 (1992), no. 2, 125–134.
[2] Noga Alon and Joel H Spencer, The probabilistic method, vol. 57, Wiley-Interscience, 2004.
[3] O.V. Borodin, Criterion of chromaticity of a degree prescription, Abstracts of IV All-Union Conf. on Th. Cybernetics, 1977, pp. 127–128.
[4] R.L. Brooks, On colouring the nodes of a network, Mathematical Proceedings of the Cambridge Philosophical Society, vol. 37, Cambridge Univ Press, 1941, pp. 194–197.
[5] G.A. Dirac, Note on the colouring of graphs, Mathematische Zeitschrift 54 (1951), no. 4, 347–353.
[6] _____, A theorem of R.L. Brooks and a conjecture of H. Hadwiger, Proceedings of the London Mathematical Society 3 (1957), no. 1, 161–195.
[7] P. Erdős, A.L. Rubin, and H. Taylor, Choosability in graphs, Proc. West Coast Conf. on Combinatorics, Graph Theory and Computing, Congressus Numerantium, vol. 26, 1979, pp. 125–157.
[8] T. Gallai, Kritische graphen i., Math. Inst. Hungar. Acad. Sci 8 (1963), 165–192 (in German).
[9] J. Hladký, D. Král, and U. Schauz, Brooks’ Theorem via the Alon-Tarsi Theorem, Discrete Mathematics (2010).
[10] H.A. Kierstead and A.V. Kostochka, Ore-type versions of Brooks’ theorem, Journal of Combinatorial Theory, Series B 99 (2009), no. 2, 298–305.
[11] A. Kostochka and M. Yancey, Ore’s Conjecture for $k = 4$ and Grötzsch Theorem, arXiv:1209.1173 (2012).
[12] ——, Ore’s Conjecture on color-critical graphs is almost true, arXiv:1209.1050 (2012).

[13] Alexandr V Kostochka, A modification of a catlin’s algorithm, Methods and Programs of Solutions Optimization Problems on Graphs and Networks 2 (1982), 75–79.

[14] A.V. Kostochka, L. Rabern, and M. Stiebitz, Graphs with chromatic number close to maximum degree, Discrete Mathematics 312 (2012), no. 6, 1273–1281.

[15] A.V. Kostochka and M. Stiebitz, A new lower bound on the number of edges in colour-critical graphs and hypergraphs, Journal of Combinatorial Theory, Series B 87 (2003), no. 2, 374–402.

[16] M. Krivelevich, On the minimal number of edges in color-critical graphs, Combinatorica 17 (1997), no. 3, 401–426.

[17] L. Rabern, ∆-critical graphs with small high vertex cliques, Journal of Combinatorial Theory, Series B 102 (2012), no. 1, 126–130.

[18] ——, Partitioning and coloring with degree constraints, Arxiv preprint arXiv:1202.5855 (2012).

[19] ——, Destroying non-complete regular components in graph partitions, J. Graph Theory (In Press).

[20] A. Riasat and U. Schauz, Critically paintable, choosable or colorable graphs, Discrete Mathematics 312 (2012), no. 22, 3373–3383.

[21] U. Schauz, Mr. Paint and Mrs. Correct, The Electronic Journal of Combinatorics 16 (2009), no. 1, R77.

[22] ——, Flexible color lists in Alon and Tarsis theorem, and time scheduling with unreliable participants, the electronic journal of combinatorics 17 (2010), no. 1, R13.

[23] V.G. Vizing, Vertex coloring with given colors, Metody Diskretn. Anal. 29 (1976), 3–10 (in Russian).

[24] X. Zhu, On-line list colouring of graphs, The Electronic Journal of Combinatorics 16 (2009), no. 1, R127.