THE COMPETITION NUMBERS OF HAMMING GRAPHS
WITH DIAMETER AT MOST THREE

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ABSTRACT. The competition graph of a digraph $D$ is a graph which has the same vertex set as $D$ and has an edge between $x$ and $y$ if and only if there exists a vertex $v$ in $D$ such that $(x, v)$ and $(y, v)$ are arcs of $D$. For any graph $G$, $G$ together with sufficiently many isolated vertices is the competition graph of some acyclic digraph. The competition number $k(G)$ of a graph $G$ is defined to be the smallest number of such isolated vertices. In general, it is hard to compute the competition number $k(G)$ for a graph $G$ and it has been one of important research problems in the study of competition graphs. In this paper, we compute the competition numbers of Hamming graphs with diameter at most three.

1. Introduction and Main Results

The notion of a competition graph was introduced by Cohen [2] in connection with a problem in ecology (see also [3]). The competition graph $C(D)$ of a digraph $D$ is the (simple undirected) graph which has the same vertex set as $D$ and has an edge between vertices $u$ and $v$ if and only if there is a vertex $x$ in $D$ such that $(u, x)$ and $(v, x)$ are arcs of $D$. For any graph $G$, $G$ together with sufficiently many isolated vertices is the competition graph of an acyclic digraph. Roberts [15] defined the competition number $k(G)$ of a graph $G$ to be the smallest number $k$ such that $G$ together with $k$ isolated vertices is the competition graph of an acyclic digraph. Opsut [12] showed that the computation of the competition number of a graph is an NP-hard problem. In the study of competition graphs, it has been one of important research problems to compute the competition numbers for various graph classes (see [1], [4], [5], [6], [7], [9], [11], [15], [17] for graphs whose competition numbers are known). For some special graph families, we have explicit formulae for computing competition numbers. For example, if $G$ is a chordal graph without isolated vertices

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then \( k(G) = 1 \), and if \( G \) is a nontrivial triangle-free connected graph then
\[
k(G) = |E(G)| - |V(G)| + 2 \quad (\text{see} \ [15]).
\]

Recently, it has been concerned to find the competition numbers of interesting graph families used in many areas of mathematics and computer science (see \([8, 10, 13, 14, 16]\)). Hamming graphs are known as an interesting graph family in connection with error-correcting codes, association schemes, and several branches of mathematics. For a positive integer \( q \), we denote the \( q \)-set \( \{1, 2, \ldots, q\} \) by \([q]\). Also we denote the set of \( n \)-tuples over \([q]\) by \([q]^n\).

For positive integers \( n \) and \( q \), the Hamming graph \( H(n, q) \) is the graph which has the vertex set \([q]^n\) and in which two vertices \( x = (x_1, x_2, \ldots, x_n) \) and \( y = (y_1, y_2, \ldots, y_n) \) are adjacent if \( d_H(x, y) = 1 \), where \( d_H : [q]^n \times [q]^n \to \mathbb{Z} \) is the Hamming distance defined by
\[
d_H(x, y) := |\{i \in [n] \mid x_i \neq y_i\}|.
\]

Note that the diameter of the Hamming graph \( H(n, q) \) is equal to \( n \) if \( q \geq 2 \). Since the Hamming graph \( H(n, q) \) is an \( n(q - 1) \)-regular graph with \( q^n \) vertices, it follows that the number of edges of the Hamming graph \( H(n, q) \) is equal to \( \frac{1}{2}n(n - 1)q^n \).

In this paper, we study the competition numbers of Hamming graphs. If \( q = 1 \), then \( H(n, 1) \) is \( K_1 \) and so the following holds:

**Proposition 1.** For \( n \geq 1 \), we have \( k(H(n, 1)) = 0 \).

If \( q = 2 \), then since \( H(n, 2) \) triangle-free and connected, we have
\[
k(H(n, 2)) = |E(H(n, 2))| - |V(H(n, 2))| + 2 = n2^{n-1} - 2^n + 2 = (n - 2)2^{n-1} + 2.
\]

**Proposition 2.** For \( n \geq 1 \), we have \( k(H(n, 2)) = (n - 2)2^{n-1} + 2 \).

If \( n = 1 \), then \( H(1, q) \) is the complete graph \( K_q \) with \( q \) vertices and so the following holds:

**Proposition 3.** For \( q \geq 2 \), we have \( k(H(1, q)) = 1 \).

However, in general, it is not easy to compute \( k(H(n, q)) \). In this paper, we give the exact values of \( k(H(2, q)) \) and \( k(H(3, q)) \). Our main results are the following:

**Theorem 4.** For \( q \geq 2 \), we have \( k(H(2, q)) = 2 \).

**Theorem 5.** For \( q \geq 3 \), we have \( k(H(3, q)) = 6 \).

We use the following notation and terminology in this paper. For a digraph \( D \), a sequence \( v_1, v_2, \ldots, v_n \) of the vertices of \( D \) is called an acyclic ordering of \( D \) if \((v_i, v_j) \in A(D) \) implies \( i > j \). It is well-known that a digraph \( D \) is acyclic if and only if there exists an acyclic ordering of \( D \). For a digraph \( D \) and a vertex \( v \) of \( D \), we define the out-neighborhood \( N^+_D(v) \) of \( v \) in \( D \) to be the set
\{w \in V(D) \mid (v, w) \in A(D)\}$, and the in-neighborhood $N_D^-(v)$ of $v$ in $D$ to be the set \{w \in V(D) \mid (v, w) \in A(D)\}. A vertex in the out-neighborhood $N_D^+(v)$ of a vertex $v$ in a digraph $D$ is called a prey of $v$ in $D$. For a graph $G$ and a vertex $v$ of $G$, we define the open neighborhood $N_G(v)$ of $v$ in $G$ to be the set \{u \in V(G) \mid uv \in E(G)\}, and the closed neighborhood $N_G[v]$ of $v$ in $G$ to be the set $N_G(v) \cup \{v\}. We denote the subgraph of $G$ induced by $N_G(v)$ (resp. $N_G[v]$) by the same symbol $N_G[v]$ (resp. $N_G[v]$).

For a clique $S$ of a graph $G$ and an edge $e$ of $G$, we say $e$ is covered by $S$ if both of the endpoints of $e$ are contained in $S$. An edge clique cover of a graph $G$ is a family of cliques of $G$ such that each edge of $G$ is covered by some clique in the family. The edge clique cover number $\theta_E(G)$ of a graph $G$ is the minimum size of an edge clique cover of $G$. An edge clique cover of $G$ is called a minimum edge clique cover of $G$ if its size is equal to $\theta_E(G)$. A vertex clique cover of a graph $G$ is a family of cliques of $G$ such that each vertex of $G$ is contained in some clique in the family. The smallest size of a vertex clique cover of $G$ is called the vertex clique cover number, and is denoted by $\theta_V(G)$.

We denote a path with $n$ vertices by $P_n$, a cycle with $n$ vertices by $C_n$, a graph with $n$ vertices and no edges by $I_n$, and a complete multipartite graph by $K_{n_1, \ldots, n_m}$.

2. Proofs of Theorems 4 and 5

2.1. Cliques in a Hamming graph

Let $\pi_j : [q]^n \to [q]^{n-1}$ be a map defined by

$$\pi_j(x_1, x_{j-1}, x_j, x_{j+1}, \ldots, x_n) = (x_1, \ldots, x_{j-1}, x_{j+1}, \ldots, x_n).$$

For $j \in [n]$ and $p \in [q]^{n-1}$, let

$$S_j(p) := \pi_j^{-1}(p) = \{x \in [q]^n \mid \pi_j(x) = p\}.$$

Then $S_j(p)$ is a clique of $H(n, q)$ with size $q$. Let

$$\mathcal{F}(n, q) := \{S_j(p) \mid j \in [n], p \in [q]^{n-1}\}.$$

Then $\mathcal{F}(n, q)$ is the family of maximal cliques of $H(n, q)$.

Lemma 6. Let $n \geq 2$ and $q \geq 2$, and let $K$ be a clique of $H(n, q)$ with size at least 2. Then there exists a unique maximal clique $S$ of $H(n, q)$ containing $K$.

Proof. Since $\mathcal{F}(n, q)$ is the family of maximal cliques of $H(n, q)$, it is sufficient to show that there is a unique maximal clique in $\mathcal{F}(n, q)$ containing $K$.

Take a vertex $x = (x_1, x_2, \ldots, x_n) \in K$. Now we will show that there exists a unique integer $j$ such that $\pi_j(x) = \pi_j(y)$ for all vertices $y \in K \setminus \{x\}$. Take a vertex $y = (y_1, y_2, \ldots, y_n) \in K \setminus \{x\}$. Since $K$ is a clique, $x$ and $y$ are adjacent. Then there is a unique integer $j \in [n]$ such that $x_j \neq y_j$ and $\pi_j(x) = \pi_j(y)$.

Suppose that there is a vertex $z = (z_1, z_2, \ldots, z_n) \in K \setminus \{x, y\}$ such that $\pi_j(z) \neq \pi_j(x)$. Since $x$ and $z$ are adjacent, there is $j_1 \in [n]$ such that $x_{j_1} \neq z_{j_1}$ and $\pi_{j_1}(x) = \pi_{j_1}(z)$, and thus $x_{j_1} = z_{j_1}$. Since $y$ and $z$ are adjacent, there is
\(j_2 \in [n]\) with \(j_2 \neq j\) such that \(\pi_{j_2}(y) = \pi_{j_2}(z)\), and thus \(y_j = z_j\). Thus we have \(x_j = z_j = y_j\), which contradicts to the fact that \(x_j \neq y_j\). Therefore, 
\(\pi_j(z) = \pi_j(x)\).

It implies that \(j\) is the unique integer such that \(\pi_j(x) = \pi_j(y)\) for all \(y \in K\). Hence \(K\) is contained in \(S_j(\pi_j(x)) \in F(n, q)\). From the uniqueness of \(j \in [n]\) and the fact that \(\pi_j(x)\) does not depend on the choice of \(x \in K\), it follows that \(S_j(\pi_j(x))\) is the unique maximal clique containing \(K\).

**Lemma 7.** The following hold:

(a) The family \(F(n, q)\) defined by (2) is an edge clique cover of \(H(n, q)\).
(b) The edge clique cover number of \(H(n, q)\) is equal to \(nq^{n-1}\).
(c) Any minimum edge clique cover of \(H(n, q)\) consists of edge disjoint maximum cliques.

**Proof.** Since each edge is contained in a maximal clique and \(F(n, q)\) is the family of maximal cliques in \(H(n, q)\), it follows that \(F(n, q)\) is an edge clique cover of \(H(n, q)\).

Let \(E\) be a minimum edge clique cover of \(H(n, q)\), that is, \(\theta_E(H(n, q)) = |E|\). Since \(F(n, q)\) is an edge clique cover with \(|F(n, q)| = nq^{n-1}\), we have \(|E| \leq nq^{n-1}\). Now we will show that \(|E| \geq nq^{n-1}\). For a clique \(S\), let \(E(S) := \left(\frac{S}{2}\right)\).

Since \(E\) is an edge clique cover of \(H(n, q)\), it holds that

\[|E(H(n, q))| \leq \sum_{S \in E} |E(S)|,\]

and the equality holds if and only if none of two distinct cliques in \(E\) have a common edge. Since the maximum size of a clique of \(H(n, q)\) is equal to \(q\), we have \(|E(S)| \leq \left(\frac{q}{2}\right)^{n-1}\) for each \(S \in E\). Therefore,

\[\sum_{S \in E} |E(S)| \leq \left(\frac{q}{2}\right)^{n-1} \times |E|,\]

and the equality holds if and only if any element of \(E\) is a maximum clique in \(H(n, q)\). Since \(|E(H(n, q))| = \frac{1}{2}n(q - 1)q^n = \left(\frac{q}{2}\right)^{n-1} \times nq^{n-1}\), it follows from (3) and (4) that \(nq^{n-1} \leq |E|\), or \(nq^{n-1} = |E|\).

Moreover, since two equalities of (3) and (4) hold, we can conclude that any minimum edge clique cover of \(H(n, q)\) consists of edge disjoint maximum cliques. \(\square\)

**Corollary 8.** The family \(F(n, q)\) defined by (2) is a minimum edge clique cover of \(H(n, q)\).

**Proof.** It follows from the fact that \(|F(n, q)| = nq^{n-1}\) and Lemma 7. \(\square\)

### 2.2. Proof of Theorem 4

In this subsection, we give a proof of Theorem 4.
Lemma 9. Let \( n \geq 2 \) and \( q \geq 2 \). For any vertex \( x \) of \( H(n,q) \), we have \( \theta_V(N_{H(n,q)}(x)) = n \).

Proof. Take any vertex \( x \in [q]^n \) of \( H(n,q) \). Then the vertex \( x \) is adjacent to a vertex \( y \) such that \( \pi_j(x) = \pi_j(y) \) for some \( j \in [n] \). We can easily check from the definition of \( H(n,q) \) that, for any \( j \in [n] \), the set \( S_j(\pi_j(x)) := \{ y \in [q]^n | \pi_j(x) = \pi_j(y) \} \) forms a clique of \( H(n,q) \). Since \( N_{H(n,q)}(x) = \cup_{j \in [n]} S_j(\pi_j(x)) \setminus \{ x \} \), the family \( \{ S_j(\pi_j(x)) | j \in [n] \} \) is a vertex clique cover of \( N_{H(n,q)}(x) \) and so \( \theta_V(N_{H(n,q)}(x)) \leq n \).

Moreover, note that \( S_j(\pi_j(x)) \cap S_{j'}(\pi_{j'}(x)) = \{ x \} \) for \( j, j' \in [n] \) where \( j \neq j' \). Take \( y_j \in S_j(\pi_j(x)) \setminus \{ x \} \) for each \( j \in [n] \). Then \( y_1, y_2, \ldots, y_n \) are \( n \) vertices of \( N_{H(n,q)}(x) \) such that no two of them can be covered by a same clique and so \( \theta_V(N_{H(n,q)}(x)) \geq n \).

\( \square \)

Opsut showed the following lower bound for the competition number of a graph.

Theorem 10 (\cite{12}). For a graph \( G \), it holds that \( k(G) \geq \min \{ \theta_V(N_G(v)) | v \in V(G) \} \).

Corollary 11. If \( n \geq 2 \) and \( q \geq 2 \), then \( k(H(n,q)) \geq n \).

Proof. It immediately follows from Lemma 9 and Theorem 10. \( \square \)

We define a total order \( \prec \) on the set \([q]^n\) as follows. Take two distinct elements \( x = (x_1, x_2, \ldots, x_n) \) and \( y = (y_1, y_2, \ldots, y_n) \) in \([q]^n\). Then we define \( x \prec y \) if there exists \( j \in [n] \) such that \( x_i = y_i \) for \( i \leq j - 1 \) and \( x_j < y_j \). The lexicographic ordering of \([q]^n\) is the ordering \( v_1, v_2, \ldots, v_q^n \) such that \( v_1 < v_2 < \ldots < v_q^n \).

Proof of Theorem 4. By Corollary 11 it follows that \( k(H(2,q)) \geq 2 \). Now we show that \( k(H(2,q)) \leq 2 \). We define a digraph \( D \) as follows:

\[
V(D) = V(H(2,q)) \cup \{ z_1, z_2 \},
\]

\[
A(D) = \left( \bigcup_{i=2}^q \{(x, (1, i - 1)) | x \in S_1(i) \} \right) \cup \left( \bigcup_{i=2}^q \{(x, (i - 1, q)) | x \in S_2(i) \} \right)
\]

\[
\cup \{(x, z_1) | x \in S_1(1) \} \cup \{(x, z_2) | x \in S_2(1) \},
\]

where \( S_j(i) \) with \( j \in \{1, 2\} \) and \( i \in [q] \) is the clique of \( H(2,q) \) defined by (1). Since \( |N_G(v) | | v \in V(D) | N_G(v) | \geq 2 \) = \( F(2,q) \), it is easy to check that \( C(D) = H(2,q) \cup \{ z_1, z_2 \} \). In addition, the ordering obtained by adding \( z_1, z_2 \) on the head of the lexicographic ordering of \( V(H(2,q)) \) is an acyclic ordering of \( D \). To see why, take an arc \( (x, y) \in A(D) \). If \( x \in S_1(1) \) or \( x \in S_2(1) \) then \( y \) is either \( z_1 \) or \( z_2 \). If \( x \in S_1(i) \) or \( x \in S_2(i) \) for some \( 2 \leq i \in [q] \), then \( x = (l, i) \) or \( x = (i, l) \) for some \( l \in [q] \). Since \( y = (1, i - 1) \) or \( y = (i - 1, q) \), we have \( y \prec x \). Therefore \( D \) is acyclic. Hence we have \( k(H(2,q)) \leq 2 \). \( \square \)
2.3. Proof of Theorem 5

In this subsection, we give a proof of Theorem 5.

2.3.1. \( k(H(3, q)) \geq 6 \) \( (q \geq 3) \)

First, we improve the lower bound for the competition number of a Hamming graph \( H(n, q) \) given in Corollary 11 in the case where \( n \geq 3 \) and \( q \geq 3 \).

**Lemma 12.** Let \( n \geq 2 \) and \( q \geq 2 \). Let \( D \) be an acyclic digraph such that \( C(D) = H(n, q) \cup I_k \) with \( I_k = \{ z_1, z_2, \ldots, z_k \} \). Let \( z_1, z_2, \ldots, z_k, v_1, v_2, \ldots, v_q \) be an acyclic ordering of \( D \). Let \( U_i := \{ v_1, \ldots, v_i \} \) for \( i = 1, \ldots, q^\prime \).

Then

\[
|\{ S \in \mathcal{F}(n, q) \mid S \cap U_i \neq \emptyset \}| \leq k + i - 1.
\]

**Proof.** Let

\[
\mathcal{N}(D) := \{ N_D^-(x) \mid x \in V(D), |N_D^-(x)| \geq 2 \},
\]

\[
\mathcal{S}_i := \{ N_D^-(x) \mid x \in U_{i-1} \cup I_k, |N_D^-(x)| \geq 2 \},
\]

\[
\mathcal{K}_i := \{ K \in \mathcal{N}(D) \mid K \cap U_i \neq \emptyset \}.
\]

Since \( D \) is acyclic, it holds that \( \mathcal{K}_i = \{ K \in \mathcal{S}_i \mid K \cap U_i \neq \emptyset \} \). Since \( |\mathcal{S}_i| \leq k + i - 1 \), it follows that

\[
|\mathcal{K}_i| = |\{ K \in \mathcal{S}_i \mid K \cap U_i \neq \emptyset \}| \leq |\mathcal{S}_i| \leq k + i - 1.
\]

For each \( K \in \mathcal{K}_i \), there exists a unique element in \( \mathcal{F}(n, q) \) containing \( K \) by Lemma 6, we denote it by \( S_K \). From (5), it remains to show that

\[
|\{ S \in \mathcal{F}(n, q) \mid S \cap U_i \neq \emptyset \}| \leq |\mathcal{K}_i|.
\]

Take \( S \in \mathcal{F}(n, q) \) such that \( S \cap U_i \neq \emptyset \). Then there exists a vertex \( x \in S \cap U_i \).

Since \( q \geq 2 \), there exists a vertex \( y \in S \setminus \{ x \} \). Since \( C(D) = H(n, q) \cup I_k \) and the vertices \( x \) and \( y \) are adjacent, there is a common prey \( u \) of \( x \) and \( y \) in \( D \).

Then \( x \in N_D^-(u) \cap U_i \) and so \( N_D^-(u) \in \mathcal{K}_i \). Since \( N_D^-(u) \) contains \( x \) and \( y \), \( S_{N_D^-(u)} \) is a maximal clique containing \( x \) and \( y \). Then both \( S \) and \( S_{N_D^-(u)} \) are maximal cliques containing \( x \) and \( y \). By Lemma 6, we have \( S = S_{N_D^-(u)} \), which implies that \( S \in \{ S_K \mid K \in \mathcal{K}_i \} \).

It follows that

\[
\{ S \in \mathcal{F}(n, q) \mid S \cap U_i \neq \emptyset \} \subseteq \{ S_K \mid K \in \mathcal{K}_i \},
\]

and together with \( |\{ S_K \mid K \in \mathcal{K}_i \}| \leq |\mathcal{K}_i|, \) (5) holds. Hence, the lemma holds. \( \square \)

**Lemma 13.** For \( n \geq 3 \) and \( q \geq 3 \), we have \( k(H(n, q)) \geq 3n - 4 \).

**Proof.** Let \( k \) be the competition number of \( H(n, q) \) and let \( D \) be an acyclic digraph such that \( C(D) = H(n, q) \cup I_k \) with \( I_k = \{ z_1, z_2, \ldots, z_k \} \). Let \( z_1, z_2, \ldots, z_k, v_1, v_2, \ldots, v_q \) be an acyclic ordering of \( D \). Let \( U_3 := \{ v_1, v_2, v_3 \} \).

By Lemma 12, it holds that

\[
|\{ S \in \mathcal{F}(n, q) \mid S \cap U_3 \neq \emptyset \}| \leq k + 2.
\]

(7)
In addition, it holds that \(|\{S \in F(n, q) \mid S \cap U_3 \neq \emptyset\}| \geq 3n - 2\) whose proof will be shown in next paragraph. Therefore, we have \(3n - 2 \leq k + 2\), or \(k \geq 3n - 4\).

Now it remains to show that \(|\{S \in F(n, q) \mid S \cap U_3 \neq \emptyset\}| \geq 3n - 2\). Consider the subgraph of \(H(n, q)\) induced by \(U_3\), say \(H\). Then \(H\) is isomorphic to one of the following:

(i) \(K_3\)  (ii) \(P_3\)  (iii) \(P_2 \cup I_1\)  (iv) \(I_3\).

**Case (i) \(H \cong K_3\):** By Lemma \ref{lem:K3}, \(U_3\) is contained in exactly one maximal clique. Without loss of generality, we may assume that \(U_3\) is contained in \(S_1((1, \ldots, 1))\), and so we may also assume that

\[
U_3 = \{(1, 1, \ldots, 1), (2, 1, \ldots, 1), (3, 1, \ldots, 1)\}.
\]

Then the family \(\{S \in F(n, q) \mid S \cap U_3 \neq \emptyset\}\) consists of the following \(3n - 2\) elements:

\[
S_1((1, \ldots, 1)), \quad S_j((i, 1, \ldots, 1)) (i \in \{1, 2, 3\}, \ j \in [n] \setminus \{1\}).
\]

**Case (ii) \(H \cong P_3\):** Without loss of generality, we may assume that

\[
U_3 = \{(1, \ldots, 1), (2, 1, \ldots, 1), (1, 2, 1, \ldots, 1)\}.
\]

Then the family \(\{S \in F(n, q) \mid S \cap U_3 \neq \emptyset\}\) consists of the following \(3n - 2\) elements:

\[
S_j((1, \ldots, 1)) \ (j \in [n]), \quad S_j((2, 1, \ldots, 1)) \ (j \in [n] \setminus \{1\}),
\]

\[
S_j((1, 2, 1, \ldots, 1)) \ (j \in [n] \setminus \{2\}).
\]

**Case (iii) \(H \cong P_2 \cup I_1\) or (iv) \(H \cong I_3\):** Let \(v\) be an isolated vertex of \(H\). Since the \(n\) cliques in \(F(n, q)\) containing \(v\) do not contain the other vertices of \(U_3\), it is sufficient to show that \(\{S \in F(n, q) \mid S \cap (U_3 \setminus \{v\}) \neq \emptyset\}\) has at least \(2n - 2\) elements. Since, for each vertex \(u \in U_3 \setminus \{v\}\), there are \(n\) cliques in \(F(n, q)\) containing \(u\) and there is at most one clique in \(F(n, q)\) containing the two vertices of \(U_3 \setminus \{v\}\). Thus we can conclude that \(\{S \in F(n, q) \mid S \cap (U_3 \setminus \{v\}) \neq \emptyset\}\) has at least \(2n - 1\) elements.

We complete the proof. \(\square\)

If \(n = 3\), then the above lower bound gives \(k(H(3, q)) \geq 5\). In this case, however, we can improve the bound as follows.

**Lemma 14.** For \(q \geq 3\), we have \(k(H(3, q)) \geq 6\).
Proof. By Lemma [13], we have \( k(H(3, q)) \geq 5 \). Suppose that \( k(H(3, q)) = 5 \). Then there exists an acyclic digraph \( D \) such that \( C(D) = H(3, q) \cup I_5 \) with 
\[ I_5 = \{z_1, z_2, \ldots, z_6\} \] 
Let \( z_1, z_2, \ldots, z_6, v_1, v_2, \ldots, v_9 \) be an acyclic ordering of \( D \). Let \( U_4 := \{v_1, v_2, v_3, v_4\} \). For convenience, let

\[ A_1 := \{S \in F(3, q) \mid S \cap U_4 \neq \emptyset\}, \quad A_2 := \{S \in F(3, q) \mid v_5 \in S\}. \]

Now we consider the subgraph \( G \) of \( H(3, q) \) induced by \( U_4 \). Any graph on 4 vertices is isomorphic to one of the following graphs:

\[
\begin{align*}
(i) & \quad K_4 & (ii) & \quad K_{1,1,2} & (iii) & \quad K_4 - E(P_3) & (iv) & \quad C_4 \\
(v) & \quad P_4 & (vi) & \quad K_{1,3} & (vii) & \quad K_3 \cup I_1 & (viii) & \quad K_2 \cup K_2 \\
(ix) & \quad P_3 \cup I_1 & (x) & \quad K_2 \cup I_2 & (xi) & \quad I_4
\end{align*}
\]

Since \( H(3, q) \) does not contain an induced subgraph isomorphic to \( K_{1,1,2} \), by Lemma [6], \( G \) is one of the above graphs except (ii). For each cases, the number \(|A_1|\) is given as follows:

\[
(i) \quad 9 \quad (ii) = (iii) \quad 9 \quad (iv) \quad 8 \\
(v) \quad 9 \quad (vi) \quad 9 \quad (vii) \quad 10 \quad (viii) \quad 10 \\
(ix) \quad 10 \quad (x) \quad 11 \quad (xi) \quad 12
\]

By Lemma [12], we have \(|A_1| \leq 8\). Therefore \( G \cong C_4 \) and so \(|A_1| = 8\). Since each vertex of \( H(3, q) \) is contained in exactly 3 cliques in \( F(3, q) \), \(|A_2| = 3\).

From the fact that

\[ A_1 \cup A_2 = \{S \in F(3, q) \mid S \cap (U_4 \cup \{v_5\}) \neq \emptyset\}, \]

it holds that \(|A_1 \cup A_2| \leq 9\) by Lemma [12]. Since \(|A_1| = 8\), \(|A_2| = 3\), and \(|A_1 \cup A_2| \leq 9\), we have \(|A_1 \cap A_2| = |A_1| - |A_1 \cup A_2| \geq 8 + 3 - 9 = 2\).

Take two distinct cliques \( S, S' \in A_1 \cap A_2 \). Then \( S \cap U_4 \neq \emptyset \), \( S' \cap U_4 \neq \emptyset \) and so take \( x \in S \cap U_4 \) and \( y \in S' \cap U_4 \). If \( x = y \) or \( x \) and \( y \) are adjacent, then \( S = S' \) by Lemma [6]. Therefore \( x \) and \( y \) are not adjacent. Since \( G \cong C_4 \), without loss of generality, we may assume that

\[ U_4 = \{(1, 1, 1), (1, 1, 2), (1, 2, 2), (1, 2, 1)\}, \quad x = (1, 1, 1), \quad y = (1, 2, 2). \]

Since \( x \) and \( v_5 \) are adjacent, one of the following holds:

\[ \pi_1(x) = (1, 1) = \pi_1(v_5), \quad \pi_2(x) = (1, 1) = \pi_2(v_5), \]

\[ \pi_3(x) = (1, 1) = \pi_3(v_5). \]

Since \( y \) and \( v_5 \) are adjacent, one of the following holds:

\[ \pi_1(y) = (2, 2) = \pi_1(v_5), \quad \pi_2(y) = (1, 2) = \pi_2(v_5), \]

\[ \pi_3(y) = (1, 2) = \pi_3(v_5). \]

However, it is impossible that \( v_5 \) satisfies both one of (8) and one of (9) since \( v_5 \notin U_4 \). We reach a contradiction. Hence we conclude \( k(H(3, q)) \geq 6 \). □
\[2.3.2. \quad k(H(3,q)) \leq 6 \quad (q \geq 3)\]

Next, we show the upper bound \(k(H(3,q)) \leq 6\) for \(q \geq 3\). To show the upper bound, we introduce a graph \(K_{q_1} \square K_{q_2} \square K_{q_3}\) as an extension of \(H(3,q)\).

For graphs \(G\) and \(H\), the Cartesian product \(G \square H\) of \(G\) and \(H\) is the graph which has the vertex set \(V(G) \times V(H)\) and has an edge between two vertices \((g,h)\) and \((g',h')\) if and only if \(gg' \in E(G)\) and \(h = h'\), or \(g = g'\) and \(hh' \in E(H)\). Note that the Cartesian product of \(n\) complete graphs \(K_q\) of size \(q\) is the Hamming graph \(H(n,q)\).

Let \(q_1,q_2,q_3 \geq 2\) be integers and we consider the graph \(K_{q_1} \square K_{q_2} \square K_{q_3}\). Define

\[
\pi_1: [q_1] \times [q_2] \times [q_3] \rightarrow [q_2] \times [q_3], \quad (x_1,x_2,x_3) \mapsto (x_2,x_3),
\]

\[
\pi_2: [q_1] \times [q_2] \times [q_3] \rightarrow [q_1] \times [q_3], \quad (x_1,x_2,x_3) \mapsto (x_1,x_3),
\]

\[
\pi_3: [q_1] \times [q_2] \times [q_3] \rightarrow [q_1] \times [q_2], \quad (x_1,x_2,x_3) \mapsto (x_1,x_2).
\]

For \(p_1 \in [q_1], p_2 \in [q_2]\), and \(p_3 \in [q_3]\), let

\[
S_1((p_2,p_3)) := \{ x \in [q_1] \times [q_2] \times [q_3] \mid \pi_1(x) = (p_2,p_3) \},
\]

\[
S_2((p_1,p_3)) := \{ x \in [q_1] \times [q_2] \times [q_3] \mid \pi_2(x) = (p_1,p_3) \},
\]

\[
S_3((p_1,p_2)) := \{ x \in [q_1] \times [q_2] \times [q_3] \mid \pi_3(x) = (p_1,p_2) \}.
\]

Note that \(S_1((p_2,p_3)), S_2((p_1,p_3)), S_3((p_1,p_2))\) are maximal cliques of \(K_{q_1} \square K_{q_2} \square K_{q_3}\). We denote the set of all maximal cliques \(S_1((p_2,p_3)), S_2((p_1,p_3))\) and \(S_3((p_1,p_2))\) by \(\mathcal{F}_{(q_1,q_2,q_3)}\). Then \(\mathcal{F}_{(q_1,q_2,q_3)}\) is an edge clique cover of \(K_{q_1} \square K_{q_2} \square K_{q_3}\).

**Lemma 15.** For \(q_1,q_2,q_3 \geq 2\), there exists an acyclic digraph \(D\) such that \(C(D) = (K_{q_1} \square K_{q_2} \square K_{q_3}) \cup I_6\) and

\[
\{N_D^-(v) \mid v \in V(D), |N_D^-(v)| \geq 2\} = \mathcal{F}_{(q_1,q_2,q_3)}.
\]

Consequently, we have \(k(K_{q_1} \square K_{q_2} \square K_{q_3}) \leq 6\).

**Proof.** For any digraph \(D\), we define \(N(D) := \{ N_D^-(v) \mid v \in V(D), |N_D^-(v)| \geq 2\} \). We prove the lemma by induction on \(m = q_1 + q_2 + q_3\). Since \(q_1,q_2,q_3 \geq 2\), we have \(m \geq 6\). Suppose \(m = 6\), i.e., \(q_1 = q_2 = q_3 = 2\). Note that \(K_2 \square K_2 \square K_2 = H(3,2)\). Since \(H(3,2)\) is a triangle-free graph, there exists an acyclic digraph \(D\) such that \(C(D) = H(3,2) \cup I_6\) and \(N_D^-(v)\) is either the empty set or a maximum clique in \(H(3,2)\) for each vertex \(v \in V(D)\) (see Figure 4 for an illustration of such a digraph). Thus the statement is true for \(m = 6\).

Suppose that the statement is true for \(m = q_1 + q_2 + q_3\) where \(m \geq 6\). Consider a graph \(K_{q_1} \square K_{q_2} \square K_{q_3}\) such that \(m + 1 = q_1 + q_2 + q_3\). Since \(q_1 + q_2 + q_3 \geq 6\), at least one of \(q_1, q_2,\) or \(q_3\) is greater than 2. Without loss of generality, we may assume that \(q_1 > 2\). Now we consider the graph \(K_{q_1-1} \square K_{q_2} \square K_{q_3}\), which is a subgraph of \(K_{q_1} \square K_{q_2} \square K_{q_3}\). Then by the induction hypothesis, there exists an acyclic digraph \(D_0\) such that \(C(D_0) = (K_{q_1-1} \square K_{q_2} \square K_{q_3}) \cup I_6\)
Figure 1. The Hamming graph $H(3, 2)$ and an acyclic digraph $D$ satisfying $C(D) = H(3, 2) \cup I_6$.

Let $v_1, v_2, \ldots, v_{q_1q_2q_3-2q_3+6}$ be an acyclic ordering of $D_0$. For convenience, let $w_1 := v_{q_1q_2q_3-2q_3+5}$ and $w_2 := v_{q_1q_2q_3-2q_3+6}$.

Let $H^*$ be the subgraph of $K_{q_1} \Box K_{q_2} \Box K_{q_3}$ induced by

$$V^* := V(K_{q_1} \Box K_{q_2} \Box K_{q_3}) - V(K_{q_1-1} \Box K_{q_2} \Box K_{q_3}) = \{ (q_1, p_2, p_3) \mid p_2 \in [q_2], p_3 \in [q_3] \}.$$

Now we define a digraph $D_1$ as follows:

$$V(D_1) = V^* \cup \{ w_1, w_2 \},$$

$$A(D_1) = \bigcup_{i=2}^{q_3} \{ (x, (q_1, 1, i-1)) \mid x \in S_2((q_1, i)) \}$$

$$\cup \bigcup_{i=2}^{q_2} \{ (x, (q_1, i-1, q)) \mid x \in S_3((q_1, i)) \}$$

$$\cup \{ (x, w_1) \mid x \in S_2((q_1, 1)) \} \cup \{ (x, w_2) \mid x \in S_3((q_1, 1)) \}.$$

The ordering obtained by adding $w_1, w_2$ on the head of the lexicographic ordering of $V^*$ is an acyclic ordering of $D_1$, and let $w_1, w_2, \ldots, w_{q_1q_2q_3+2}$ be the ordering. In addition,

$$\mathcal{N}(D_1) = \{ S_2((q_1, i)) \mid i \in [q_3] \} \cup \{ S_3((q_1, i')) \mid i' \in [q_2] \}.$$

Therefore $D_1$ is an acyclic digraph such that $C(D_1) = H^* \cup \{ w_1, w_2 \}$.

Note that, for $(p_2, p_3) \in [q_2] \times [q_3]$, the clique in $K_{q_1} \Box K_{q_2} \Box K_{q_3}$ obtained by deleting the vertex $(q_1, p_2, p_3)$ from an element $S_1((p_2, p_3))$ of $K_{q_1} \Box K_{q_2} \Box K_{q_3}$.
is a maximal clique of $K_{q_1-1} □ K_{q_2} □ K_{q_3}$. Then by (10), for each $(p_2, p_3) \in [q_2] \times [q_3]$, the set \( \{ v \in V(D_0) \mid N_{D_0}^{-}(v) = S_1((p_2, p_3)) \setminus \{(q_1, p_2, p_3)\} \} \) is not empty, and so we take an element $y_{(p_2, p_3)}$ of this set.

Now we define a digraph $D$ as follows:

\[
V(D) = V(K_{q_1} □ K_{q_2} □ K_{q_3}) \cup I_6
\]
\[
A(D) = A(D_0) \cup A(D_1) \cup \{ ((q_1, p_2, p_3), y_{(p_2, p_3)}) \mid (p_2, p_3) \in [q_2] \times [q_3] \}
\]

Note that since $N_{D_0}(w_1) = N_{D_0}(w_2) = \emptyset$, (12)

\[
\{ N_D(w_1), N_D(w_2) \} = \{ S_2((q_1, 1)), S_3((q_1, 1)) \}.
\]

By the definition of $D$ and (10), (11), and (12), we can conclude that $N(D) = F_{(q_1, q_2, q_3)}$ and so $E(C(D)) = E(K_{q_1} □ K_{q_2} □ K_{q_3})$. Thus, $C(D) = (K_{q_1} □ K_{q_2} □ K_{q_3}) \cup I_6$. Then the ordering

\[
v_1, v_2, \ldots, v_{q_1 q_2 q_3 - q_2 q_3 + 6}, w_3, w_4, \ldots, w_{q_2 q_3 + 2}
\]

of the vertices of $D$ is an acyclic ordering. To see this, take an arc $a = (x, y) \in A(D)$. If $a \in A(D_0) \cup A(D_1)$, then $y$ appears before $x$ in (13), since $D_0$ and $D_1$ are acyclic. If $a \not\in A(D_0) \cup A(D_1)$ then $x \in \{ w_3, w_4, \ldots, w_{q_2 q_3 + 2} \}$ and $y \in \{ v_1, v_2, \ldots, v_{q_1 q_2 q_3 - q_2 q_3 + 6} \}$, thus $y$ appears before $x$ in (13). Thus the digraph $D$ is acyclic. Hence the lemma holds. \( \square \)

**Proof of Theorem 5** By Lemma 14 we have $k(H(3, q)) \geq 6$ for $q \geq 3$. By Lemma 15 we have $k(H(3, q)) = k(K_q □ K_q □ K_q) \leq 6$ for $q \geq 3$. Hence Theorem 5 holds. \( \square \)

### 3. Concluding Remarks

In this paper, we gave the exact values of the competition numbers of Hamming graphs with diameter 2 or 3.

We conclude this paper with leaving the following questions for further study:

- What is the competition number of a Hamming graph $H(4, q)$ with diameter 4 for $q \geq 3$?
- What is the competition number of a ternary Hamming graph $H(n, 3)$ for $n \geq 4$?
- Give the exact values or a good bound for the competition numbers of Hamming graphs $H(n, q)$.

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