Detection Thresholds in Very Sparse Matrix Completion

Charles Bordenave · Simon Coste · Raj Rao Nadakuditi

Received: 15 May 2020 / Revised: 4 August 2021 / Accepted: 9 December 2021 / Published online: 14 June 2022
© SFoCM 2022

Abstract
We study the matrix completion problem: an underlying $m \times n$ matrix $P$ is low rank, with incoherent singular vectors, and a random $m \times n$ matrix $A$ is equal to $P$ on a (uniformly) random subset of entries of size $dn$. All other entries of $A$ are equal to zero. The goal is to retrieve information on $P$ from the observation of $A$. Let $A_1$ be the random matrix where each entry of $A$ is multiplied by an independent $\{0, 1\}$-Bernoulli random variable with parameter $1/2$. This paper is about when, how and why the non-Hermitian eigen-spectra of the matrices $A_1(A - A_1)^*$ and $(A - A_1)^*A_1$ captures more of the relevant information about the principal component structure of $A$ than the eigen-spectra of $AA^*$ and $A^*A$. We show that the eigenvalues of the asymmetric matrices $A_1(A - A_1)^*$ and $(A - A_1)^*A_1$ with modulus greater than a detection threshold are asymptotically equal to the eigenvalues of $PP^*$ and $P^*P$ and that the associated eigenvectors are aligned as well. The central surprise is that by intentionally inducing asymmetry and additional randomness via the $A_1$ matrix, we can extract more information than if we had worked with the singular value decomposition (SVD) of $A$. The associated detection threshold is asymptotically exact and is non-universal since it explicitly depends on the element-wise distribution of the underlying matrix $P$. We show that reliable, statistically optimal but not perfect matrix recovery, via a universal data-driven algorithm, is possible above this detection threshold.
threshold using the information extracted from the asymmetric eigen-decompositions. Averaging the left and right eigenvectors provably improves estimation accuracy but not the detection threshold. Our results encompass the very sparse regime where $d$ is of order 1 where matrix completion via the SVD of $A$ fails or produces unreliable recovery. We define another variant of this asymmetric principal component analysis procedure that bypasses the randomization step and has a detection threshold that is smaller by a constant factor but with a computational cost that is larger by a polynomial factor of the number of observed entries. Both detection thresholds allow to go beyond the barrier due to the well-known information theoretical limit $d \approx \log n$ for exact matrix completion found in the literature.

**Keywords** Matrix completion · Sparse random graphs · Eigenvalues · Spectral algorithms · Non-Hermitian matrices

**Mathematics Subject Classification** 60B20 · 15B52 · 05C80 · 62R01

**Contents**

1 Introduction ............................................. 1620  
2 Main Results on Matrix Completion ................................ 1627  
3 Related Work and Extensions ................................. 1631  
4 Detailed Results: Square Matrices ......................... 1636  
5 Detailed Results: Non-Backtracking Matrix ............. 1642  
6 Detailed Results: Rectangular Matrices ................. 1647  
7 The Rank-one Case: Complete Theory .................. 1658  
8 An Algebraic Perturbation Lemma ....................... 1666  
9 Proof of Theorem 4.3 .................................... 1677  
10 Consequences of Algebraic Incoherence .............. 1683  
11 Coupling Graphs and Trees ............................ 1686  
12 Graph Functionals ..................................... 1692  
13 Near Eigenvectors: Proof of Theorem 9.2 .......... 1696  
14 Eigenwaves on Galton-Watson Trees: Proof of Theorem 12.2 ............ 1704  
15 Proof of Proposition 13.5 ................................ 1708  
16 Proof of Theorem 5.4 and Corollary 5.5 .............. 1717  
17 Proofs for the Rectangular case ....................... 1733  
18 Some Possible Improvement of the Quantitative Bounds .................. 1740  
References ............................................... 1741

1 Introduction

1.1 Setting and Overview

We start by describing the main mathematical model that will be studied in this paper. Let $m, n \geq 1$ large integers, and let $P = (P_{x,y}) \in \mathcal{M}_{m,n}(\mathbb{R})$ be a real matrix with...
singular value decomposition

\[ P = \sum_{k=1}^{r} \sigma_k \zeta_k \xi_k^*, \quad (1.1) \]

the positive numbers \( \sigma_k \) are the singular values of \( P \), and \((\zeta_1, \ldots, \zeta_r), (\xi_1, \ldots, \xi_r)\) are two orthonormal families of singular vectors. Let \( M \in \mathcal{M}_{m,n}(\mathbb{R}) \) be a random matrix whose entries are independent Bernoulli with parameter \( d/n \): for all \( x, y \), we have

\[ P(M_{x,y} = 1) = 1 - P(M_{x,y} = 0) = \frac{d}{n}. \]

The non-zeros entries of the matrix \( M \) correspond to the entries of the matrix \( P \) that are observed, the remaining ones are hidden. The observed matrix is then defined as

\[ A = \left( \frac{n}{d} \right) P \odot M, \]

where \( \odot \) denotes the Hadamard entry-wise product of two matrices. The normalization is chosen so that \( \mathbb{E}[A] = P \), hence \( A \) is an unbiased estimator of \( P \). The matrix \( A \) has an average of \( d \) revealed entries per row. From now on, we will concentrate on the asymptotic regime where \( m \) and \( n \) are large and have the same order, by supposing that \( \alpha = m/n \) is bounded away from 0 and \( \infty \).

The matrix completion problem aims at answering the following general question: what parts of \( P \) can be recovered from the observed entries? The literature around this problem is gigantic, see Sect. 3. Roughly speaking, it is known that under natural assumptions on \( P \) (low-rank with delocalized eigenvectors), we can recover exactly \( P \) as soon as \( d \) has order \( \log n \). Below this threshold, there exists an estimator \( \hat{P} \) whose mean square error, that is \( \text{tr}(\hat{P} - P)^2/n \), is of order \( 1/d \).

In this paper, we go beyond this last result and study the detection problem of the spectrum of \( P \). Namely, for a given singular value \( \sigma_k \) of \( P \) with corresponding unit singular vectors \( \zeta_k, \xi_k \), our goal is to address the following questions:

(i) For which values of \( d \) is it possible to design a consistent estimator of \( \sigma_k \)?
(ii) For a given \( \varepsilon > 0 \), for which values of \( d \) is it possible to design an estimator \( \hat{\zeta}_k \) of \( \zeta_k \) such that \( |\langle \zeta_k, \hat{\zeta}_k \rangle| \geq \varepsilon \) with high probability, and the same for \( \xi_k \)?
(iii) If such estimators can be found, what is the best way to use them for estimating the matrix \( P \)?

The results in the existing literature imply that (i) and (ii) are feasible when \( d \) goes to infinity. Our results prove that this is possible for \( d \) of order 1 simply by considering the \( k \)-th largest eigenvalue of an \( m \times m \) carefully chosen matrix \( X \) and its corresponding eigenvectors. We then adapt the OptShrink strategy [48] to answer (iii) and give asymptotically exact mean-square errors for the estimation of \( P \) with our technique.
1.2 Informal Statement of the Main Results

1.2.1 When $P$ is Hermitian

When $P$ is itself a Hermitian matrix, i.e., $P = P^*$, then given any fixed $d$, there is a threshold $\vartheta$, intrinsic to the matrix $P$ and to $d$, such that the following holds with high probability when $n$ is large.

(i) Each eigenvalue $\sigma_i > \vartheta$ gives rise to an eigenvalue $\lambda_i$ of $A$ close to $\sigma_i$. The rest of the eigenvalues of $A$ are quarantined in the disc $D(0, \vartheta)$.

(ii) Moreover, if $\chi_i$ is a unit right-eigenvector of $A$ associated with the eigenvalue $\lambda_i \sim \sigma_i$ above the threshold $\vartheta$, then the scalar product between $\chi_i$ and the singular vector $\zeta_i$ has an explicit non-vanishing limit.

1.2.2 When $P$ is Rectangular

When $P$ is rectangular, we adopt a downsampling/dilation strategy. Let $Z \in \mathbb{M}_{m,n}(\mathbb{R})$ be an auxiliary random matrix whose entries are independent Bernoulli with parameter $1/2$. We set $A_1 = Z \circ A$ and $A_2 = A - A_1$ and we define

\[ X = A_1 A_2^* \quad \quad \quad Y = A_1^* A_2. \] (1.2)

These are square matrices, with respective sizes $m \times m$ and $n \times n$. They are not Hermitian, their eigenvalues are complex numbers. Then, given any fixed $d$, there is a threshold $\vartheta$, intrinsic to the matrix $P$ and to $d$, such that the following holds with high probability when $n$ is large.

(i) Each singular value $\sigma_i > \vartheta$ gives rise to an eigenvalue $\nu_i$ of $X$ close to $\sigma_i^2$. The rest of the eigenvalues of $X$ are quarantined in the disc $D(0, \vartheta)$.

(ii) Moreover, if $\chi_i$ is a unit right-eigenvector of $X$ associated with the eigenvalue $\nu_i \sim \sigma_i^2$ above the threshold $\vartheta$, then the scalar product between $\chi_i$ and the singular vector $\zeta_i$ has an explicit non-vanishing limit.

Similar results hold for $Y$ and $\xi_i$.

1.2.3 Statistical Reconstruction

In both cases ($P$ Hermitian or rectangular), all the theoretical quantities can efficiently be estimated from the observation of $P \circ M$, leading to asymptotically efficient data-driven estimators even in the sparsest regime where $d$ is fixed. This threshold will be essentially the analog of the Kesten-Stigum threshold well known in the field of community detection [14]. When all the singular values of $P$ are above the threshold $\vartheta$, we can thus use the eigenvalues and eigenvectors of $X$ and $Y$ to craft an estimator $\hat{P}$ of $P$ which is correlated with $P$. 

\[ \hat{P} \]

\[ \text{Springer} \]
1.2.4 Strengths of this Method

This spectral method does not need complex manipulations, does not require trimming, uses all the available information, and only needs computing the top eigenvalues and eigenvectors of $m \times m$ or $n \times n$ matrices, hence is computationally efficient. On the other hand, spectral algorithms are generally known to be less robust than other methods, and the detection threshold is 1.44 times higher than the so-called non-backtracking threshold studied in our Sect. 5, the latter gives better results but requires a higher computational cost.
1.3 A Toy Model for Illustrating our Results

Our results are part of a new and promising philosophy, namely that in many problems eigenvalues of non-symmetric matrices can perform better than eigenvalues of symmetric matrices, notably initiated in [23]. We will highlight this statement with a toy model and a preview of our numerical experiments from Sect. 7.1.

Suppose that the underlying unknown matrix $P$ has rank 1 and is symmetric, say $P = \varphi \varphi^*$; each entry is observed with probability $d/n$, resulting in a normalized observed matrix $A$. Since the resulting observations are not Hermitian but the underlying problem is Hermitian, one might be tempted to say that if the entry $(i, j)$ of $P$ is observed then so is the entry $(j, i)$: thus, setting $A_{j, i} = A_{i, j}$ seems like a reasonable processing, before applying a PCA to the ‘symmetrized observations’ $\tilde{A}$: that is, compute the highest eigenvalue $\lambda_1(\tilde{A})$ and eigenvector $u$ of $\tilde{A}$, and estimate $P$ with $\lambda_1(\tilde{A})uu^*$. We will see that this method yields poor results: instead, it is much more efficient to keep the observations $A$ non-Hermitian, to compute the top eigenvalue and the left/right eigenvectors $u, v$, and to estimate $P$ with $\lambda_1(A)uv^*$.

In general, when $P$ is a rectangular matrix, we want to use the same strategy. Since the observed matrix is rectangular, there are no eigenvalues, and one might be tempted to directly use its singular values, which are the square roots of $AA^*$. Here again, $AA^*$ is Hermitian, and this method yields poor results; instead, it is more efficient to downsample $A$ into $A_1$ and use the eigenvalues of $A_1(A - A_1)^*$, which is a non-Hermitian analog of $AA^*$. To study this model, we simply reduce to the Hermitian case described above by a ‘symmetric dilation trick’ described above and in Fig. 2; but the core strength of our results still come from the asymmetric-eigendecomposition idea, not from the dilation trick.

The numerical experiments in Fig. 1 illustrate these phenomena. The first column of Fig. 1 displays the $30 \times 10$ upper-left sub-matrices of $2000 \times 3000$ rank one matrices modeled as in (1.1). In Fig. 1a the singular vectors are normally distributed whereas in Fig. 1b they are drawn from the hyperbolic secant distributions. This matrix was very sparsely sampled ($d = 9.7$ and $d = 22.6$ for the Gaussian and Hyperbolic setting, respectively). The second, third and fourth columns of Fig. 1 show reconstructions obtained using the SVD (with the missing entries replaced by zeros), the eigen-spectra of the asymmetric matrix as described above (which averages the left and right eigenvectors of the $X$ and $Y$ matrices to produce an estimate of the left and right singular vectors, respectively) and also an additional refinement that will be defined later, using the right eigenvector of the weighted non-backtracking matrix.

The SVD reconstruction fails completely, whereas the two asymmetric methods described succeed and produce reliable, even if imperfect, estimates. To summarize the emergent philosophy: we succeed in finding (symmetric) structure in a regime where a symmetric eigen-decomposition fails by inducing, via randomization, asymmetry and viewing the symmetric problem through an asymmetric eigen-decomposition lens.
1.4 Sneak Peek: Improved Matrix Recovery Using Asymmetric Eigen-Methods

The reason why eigenvalues of symmetric matrices or singular values are inefficient in reconstructing a signal $P$ from a sparse set of observations mainly comes from the sensitivity of eigenvalues to heterogeneity.

When sampling each entry with a probability $d/n$, the number of revealed entries on each line will be highly heterogeneous. Some lines will have zero revealed entries, while some lines will have a maximum of $\log n$ entries revealed. The number of revealed entries is not concentrated, unlike denser case where for instance the unmasking probability is $\omega(\log(n)/n)$.

Due to the variational characterizations of eigenvalues of Hermitian matrices / singular values of general matrices, the highest eigen or singular values of $A$ will tend to be almost entirely determined by the few rows and columns with many revealed entries, and the corresponding eigen/singular vectors will be localized on the $O(\log n)$ revealed entries of these rows, and will thus only catch information on a very tiny part of $P$. This is why the celebrated paper [37] introduced a spectral regularization procedure consisting in trimming all those high degrees, sometimes at the cost of throwing a lot of information.
The eigenvalues of non-normal matrices do not suffer this localization effect, precisely because they do not obey variational inequalities. In our work, we propose to exploit this by crafting different non-Hermitian matrices whose spectrum is provably well correlated with $P$: the observation $A$ when $P$ is Hermitian, the downsampled versions $X$, $Y$ in the rectangular case, or even a variant of the non-backtracking matrix, to be defined later.

This plan (artificially introducing non-symmetry) was spectacularly successful for solving the community detection problem. In this problem, one observes a graph sampled from a generative model with planted communities (the SBM), and tries to recover the communities; in the sparse regimes, the degrees of the graph are very heterogeneous and thus the usual spectral clustering procedures do not work well, for the reason exposed above. The spectral redemption conjecture [39] stated that using the eigenvalues of a non-Hermitian matrix (the non-backtracking matrix) could succeed in detecting the communities, and it was proven true in the breakthrough paper [14]. In this paper, they describe a phase transition phenomenon well-known as the Kesten-Stigum threshold: under some condition on the connectivities inside the communities, the spectrum of the non-backtracking matrix contains informations that were not available in the spectrum of the laplacian. Regarding this phenomenon and the link between our paper and the KS threshold, we refer to the related work section. Later [23] added weight to this philosophy, in a different setting. They showed that the eigenvalues of non-Hermitian Gaussian deformations of low-rank matrices were more informative than singular values, more robust to noise heteroscedasticity.

To summarize these various findings, including the results of our paper: all other things being equal, practitioners can expect greater gains with these methods for recovering incoherent matrices whose elements have a larger kurtosis.

We hope that practitioners who encounter low-rank matrix completion problems in high-dimensional, very sparse settings, such as in reinforcement learning and computational game theory where reward/payoff matrices are often modeled this way [36, 53], can utilize these methods to find structure in severely under-sampled regimes where the failure of SVD based matrix completion methods might have been misconstrued as a by-product of a fundamental informational barrier. We shall release a numerical implementation of our methods to facilitate revised experimentation on such given-up-for-being-too-sparse data sets.

1.5 Notation and Convention

When $n$ is an integer, $[n]$ denotes the set $\{1, \ldots, n\}$. The group of permutations of $[r]$ is noted $\mathcal{S}_r$. We identify $\mathbb{R}^n$ with the set $\ell^2([n])$. Elements in $\mathbb{R}^n$ will be noted $u = (u(x))_{x \in [n]}$. We will note $|\cdot|_\infty$, $|\cdot|_p$ the usual norms on $\mathbb{R}^n$, namely

$$
|u|_\infty = \max_{x \in [n]} |u(x)| \quad |u|_p = \left( \sum_{x \in [n]} |u(x)|^p \right)^{1/p}.
$$

The Euclidean norm ($p = 2$) will simply be noted $|\cdot|$. The operator norm of the matrix $X$ is noted $\|X\|$; it is the greatest singular value of the matrix. The Frobenius norm

$$
\hat{\cdot}^{\text{Fro}} \quad \hat{\cdot}^{\text{Fro}}
$$
is noted $\|X\|_F$ and is defined by $\|X\|_F = \sqrt{\text{tr}(X^*X)}$. It is also the $L^2$-norm of the singular values.

The letter $c$ denotes a universal numerical constant. It might be used from line to line to denote different constants.

We will also make the following convention on the phase of eigenvectors. If $\psi$ and $\psi'$ are two right and left eigenvectors of a matrix associated to the same simple eigenvalue, we will also assume that their phase is chosen so that

$$
(\psi', \psi) \geq 0.
$$

(1.3)

### 2 Main Results on Matrix Completion

Let $m \leq n$ be two integers, and let $P$ be an $m \times n$ real matrix with singular value decomposition given by

$$
P = \sum_{i=1}^{r} \sigma_i \zeta_i \xi_i^*.
$$

We want to find back $P$ from the observation of its masked version $P \odot M$, where the probability of uncovering each entry is $d/n$ for a fixed $d$. What we call the observed matrix is

$$
A = \left(\frac{n}{d}\right) P \odot M,
$$

(2.1)

where the normalization with $n/d$ is here to ensure that $\mathbb{E}[A] = P$. Let $Q$ be the $m \times n$ matrix defined by $Q_{x,y} = n|P_{x,y}|^2$ and $\rho = \|Q\|$.

**Definition 2.1** The detection threshold is defined as the real number

$$
\vartheta = \max \left( \sqrt{\frac{\rho}{2d}}, \frac{\|P\|_\infty}{2d} \right).
$$

The detection rank $r_0$ is defined as the number of singular values of $P$ which are strictly greater than this threshold: $r_0 = \max\{i \in [r] : \sigma_i > \vartheta\}$.

We will note $P_0$ the matrix $P$ whose SVD is truncated at $r_0$:

$$
P_0 := \sum_{i=1}^{r_0} \sigma_i \zeta_i \xi_i^*.
$$

The main message of this paper is that one can efficiently estimate $P_0$ from the observation of $A$; before explaining the procedure, we introduce the general method and our notion of loss.
2.1 Mean Squared Error Optimal Matrix Recovery

Suppose that we have at our disposition some estimators \( \hat{\zeta}_i \) and \( \hat{\xi}_i \) of the singular vectors \( \zeta_i, \xi_i \) — we do not specify what they are for the moment. Then, we can try to estimate \( \mathbb{P}_0 \) by

\[
\sum_{i=1}^{r_0} w_i \hat{\zeta}_i \hat{\xi}_i^* 
\]

for some well-chosen numbers \( w_i \); the best choice of the coefficients \( w_i \) is given by the solution of the optimisation problem (see [48]):

\[
w_{\text{opt}} = (w_{\text{opt}}^1, \ldots, w_{\text{opt}}^{r_0}) = \arg \min_{w_1, \ldots, w_{r_0} > 0} \| \mathbb{P}_0 - \sum_{i=1}^{r_0} w_i \hat{\zeta}_i \hat{\xi}_i^* \|_F. \tag{2.2}
\]

This problem can be solved using elementary analysis, the solution being

\[
w_i^{\text{opt}} = \left( \Re \left( \sum_{j=1}^{r_0} \sigma_j \langle \zeta_j, \hat{\zeta}_i \rangle \langle \xi_j, \hat{\xi}_i \rangle \right) \right)_+, \tag{2.3}
\]

see [48], Theorem 2.1a and the proof therein. The optimal approximation of \( \mathbb{P}_0 \) given \( \hat{\zeta}_i, \hat{\xi}_j \) is thus given by

\[
\sum_{i=1}^{r_0} w_i^{\text{opt}} \hat{\zeta}_i \hat{\xi}_i^*. \tag{2.4}
\]

There is however a problem with this method: \( w_i^{\text{opt}} \) in (2.3) is not a statistic, because it depends on the quantities \( \sigma_i, \langle \zeta_j, \hat{\zeta}_i \rangle \) and \( \langle \xi_j, \hat{\xi}_i \rangle \) which are \textit{a priori} unknown. Fortunately, our method consists in finding estimators \( \hat{\zeta}_i, \hat{\xi}_j \) for which all these quantities can be consistently estimated directly from the data.

2.2 Procedure for Finding Estimators of the Eigenvectors

The starting data are the observed entries of \( \mathbb{P} \); we note \( T = \mathbb{P} \odot M \) the \( m \times n \) observed matrix with no normalization. We are going to subsample \( T \): let \( Z' \) be an \( m \times n \) matrix with i.i.d. Bernoulli \( 1/2 \) entries. We set \( C_1 = T \odot Z \) and \( C_2 = T - C_1 \), and

\[
X = \left( \frac{2n}{d} \right)^2 C_1 C_2^* \quad Y = \left( \frac{2n}{d} \right)^2 C_2^* C_1. \tag{2.5}
\]

These matrices are respectively of size \( m \times m \) and \( n \times n \). For \( i = 1, \ldots, r_0 \), let \( \chi_i, \chi_i' \) be two unit left and right eigenvectors of \( C_1 \) associated with the eigenvalues of \( C_1 \) with greatest magnitude. Similarly, for \( i = 1, \ldots, r_0 \), let \( \pi_i, \pi_i' \) be two unit left and right eigenvectors of \( C_2 \) associated with the eigenvalues of \( C_2 \) with greatest magnitude.
Algorithm 1 Statistically optimal matrix completion using asymmetric eigen-decomposition

1: **Input**: $T = P \odot M$ the masked matrix.
2: **Input**: $r_0$ = estimate of the detection rank $r_0$.
3: Generate $X, Y$ as in (2.5).
4: Compute the $r_0$ highest eigenvalues $v_i$ of $X$, the associated right-eigenvectors $\chi_i$, the left eigenvectors $\chi'_i$. Select an orientation such that $\langle \chi_i, \chi'_i \rangle \geq 0$.
5: Compute the $r_0$ highest eigenvalues $\eta_i$ of $Y$, the associated right-eigenvectors $\pi_i$, the left eigenvectors $\pi'_i$. Select an orientation such that $\langle \pi_i, \pi'_i \rangle \geq 0$.
6: for $i = 1, \ldots, r_0$ do
7: Set
\[
\hat{\zeta}_i = \frac{\chi_i + \chi'_i}{|\chi_i + \chi'_i|} \quad \hat{\xi}_i = \frac{\pi_i + \pi'_i}{|\pi_i + \pi'_i|}. \tag{2.7}
\]
8: Set
\[\hat{w}_i = \sqrt{\frac{4v_i \langle \chi_i, \chi'_i \rangle \langle \pi_i, \pi'_i \rangle}{(1 + \langle \chi_i, \chi'_i \rangle)(1 + \langle \pi_i, \pi'_i \rangle)}}.\]
9: end for
10: return $\hat{P} = \sum_{i=1}^{r_0} \hat{w}_i \hat{\zeta}_i (\hat{\xi}_i)^*.$

Our estimators are
\[
\hat{\zeta}_i = \frac{\chi_i + \chi'_i}{|\chi_i + \chi'_i|} \quad \hat{\xi}_i = \frac{\pi_i + \pi'_i}{|\pi_i + \pi'_i|}. \tag{2.6}
\]

Most of this paper consists in proving that these estimators are correlated with $\zeta_i, \xi_j$; more precisely, we compute the exact limit of quantities like $\langle \hat{\zeta}_i, \xi_j \rangle$.

Now, we must still estimate $\sigma_i$ and $\langle \hat{\zeta}_i, \hat{\zeta}_i \rangle$ and $\langle \hat{\xi}_i, \hat{\xi}_i \rangle$ to obtain an estimation $\hat{w}_i$ of $w_i^{opt}$. We will explain in all details how we do this in the future sections, but for the moment we simply give the endpoint of our theory: we obtain a remarkably simple estimator given by
\[\hat{w}_i = \sqrt{\frac{4v_i \langle \chi_i, \chi'_i \rangle \langle \pi_i, \pi'_i \rangle}{(1 + \langle \chi_i, \chi'_i \rangle)(1 + \langle \pi_i, \pi'_i \rangle)}}.\]

This can be computed directly from the knowledge of $v_i, \pi_i, \pi'_i, \chi_i, \chi'_i$ which were extracted in the preceding step.

Putting all things together, we obtain the following final estimation for $P_0$:
\[\hat{P} = \sum_{i=1}^{r_0} \hat{w}_i \hat{\zeta}_i \hat{\xi}_i.\]

The algorithmic description of this method is described in Algorithm 1 at page 8. We now turn to a detailed analysis of the error obtained when estimating $P$ with $\hat{P}$.
2.3 The Limit of the MSE

Given the estimator \( \hat{P} \) obtained with our method, the quantity \( \| P_0 - \hat{P} \|_F \) is an approximation of the optimal mean-square error we can get with our estimators. We note it MSE\(_\star\):

\[
\text{MSE}_\star = \| P_0 - \hat{P} \|_F^2.
\] (2.8)

Remarkably, we could extract from our theory the exact limiting value of this error, as a function of the problem parameters \( P \) and \( d \). The conditions under which the result is valid are that \( P \) is sufficiently low-rank and delocalized: more precisely, we ask that \( \| P \|_\infty = O(1/n) \) (delocalized entries), that \( |\zeta_i|_\infty, |\xi_i|_\infty = O(1/\sqrt{n}) \) (delocalized eigenvectors) and that \( \| P \|_F = O(\|P\|) \) (the stable numerical rank is bounded, an assumption weaker than \( r = O(1) \)). We introduce a notation:

\[
\Gamma^{\triangle}_{i,j} = \sum_{s=0}^{\infty} \frac{1}{(2\sigma_i^2 n/d)^s} \left( 1, \begin{pmatrix} 0 & Q^* \\ Q & 0 \end{pmatrix}^s \left( \zeta_i \otimes \zeta_j \right) \right)
\] (2.9)

\[
\Gamma^{\triangledown}_{i,j} = \sum_{s=0}^{\infty} \frac{1}{(2\sigma_i^2 n/d)^s} \left( 1, \begin{pmatrix} 0 & Q^* \\ Q & 0 \end{pmatrix}^s \left( 0 \otimes \xi_i \right) \right)
\] (2.10)

These quantities only depend on the spectral decomposition of \( P \) and \( d \). No simpler expressions were found in the general case.

Proposition 2.2 (Minimum square error) On an event with probability tending to 1, the mean square error MSE\(_\star\) in (2.8) obtained with the aforementioned estimators is asymptotically given by

\[
\sum_{i=1}^{m} \sigma_i^2 \left( 1 - \frac{8}{(\Gamma^{\triangle}_{i,i} + 1)(\Gamma^{\triangledown}_{i,i} + 1)} \right) + \sum_{i,j \in [r_0]} \sigma_i \sigma_j \frac{4(\Gamma^{\triangle}_{i,j} + \delta_{i,j})(\Gamma^{\triangledown}_{i,j} + \delta_{i,j})}{(\Gamma^{\triangle}_{i,i} + 1)^2(\Gamma^{\triangledown}_{j,j} + 1)^2}
\] (2.11)

in the sense that \( \frac{\text{MSE}_\star}{(2.11)} \) converges to 1 in probability.

It is remarkable that this mean-square error can be exactly identified and the limiting value can be computed in terms of the initial problem. The proof is at Sect. 17.6. In Sect. 7, we provide a comprehensive example where the computations lead to a simple expression and we check this expression with numerical simulations. Also, it is easily understood from the definition that when \( d \to \infty \), the quantity \( \Gamma^{\triangle}_{i,j} \) converges towards 1 if \( i = j \) and zero if \( i \neq j \). Consequently, from 2.11 we have

\[
\lim_{d \to \infty} \limsup_{n \to \infty} \text{MSE}_\star = 0
\]

thus ensuring that in the \( d \to \infty \) regime, recovery is almost exact.
2.4 Roadmap of the Paper and Auxiliary Results

Section 3 surveys the rich literature on sparse spectral graph theory, random matrices, principal component analysis and matrix completion. Subsequent sections are divided into (i) Detailed statements of the results (Sects. 4–6), (ii) A fully treated example (Sect. 7), and (iii) All the technical details and proofs (Sect. 8 et sequitur, starting at page 49).

The results from the preceding section will indeed follow from the simpler case where \( P \) is a square, Hermitian matrix. In this case, we do not need to form the matrices \( X \) and \( Y \) above: the observed matrix, \( A \), is itself a square matrix, hence we can directly show the aforementioned phenomenon directly on \( A \). The setup is depicted in Fig. 2a. Informally we will show that when the underlying is symmetric and low-rank and we observe missing entries as in the left panel of Fig. 2a, then the randomly asymmetric matrix formed from the original matrix has (left and right) eigenvectors that are well-aligned with the eigenvectors of the underlying low-rank symmetric matrix. The detailed statement of this result is contained in Theorem 4.3.

The results on rectangular matrices will then be obtained through a Hermitization trick, by considering the matrix

\[
\begin{pmatrix}
0 & P \\
P^* & 0
\end{pmatrix}
\] (2.12)

and applying our theorem for square matrices as depicted in Fig. 2b.

We included in Sect. 5 several results on non-backtracking matrices, when the underlying \( P \) is square. The very recent paper [54] was built on a preliminary version of the present work and generalizes this portion to the case of weighted inhomogeneous undirected graphs.

All the subsequent sections, starting with Sect. 8 at page 39, are the technical proofs of our results. A few technical results which were developed in the course of the proof might be of independent interest. Among them, we mention the perturbation results from Sect. 8, dealing with spectra of perturbations of non-normal matrices (Theorem 9.2), and also the results of Sect. 12 which include new and powerful concentration inequalities for functionals on Erdős-Rényi random graphs (see Proposition 12.3).

3 Related Work and Extensions

A first version of this paper appeared in the PhD manuscript of the second author, in 2019 [25].

3.1 Completion and Sparsification

The problem of sparse completion consists in observing a very sparse sample of elements of a general object (a matrix, a subspace) carrying some structure (low-rank, delocalized), and trying to reconstruct it. The problem of matrix completion has attracted a gigantic amount of attention from researchers in applied mathematics since
the last 15 years; the general philosophy can be grasped by a handful of seminal papers from Candès and Tao [20] and Candès and Recht [19], Keshavan Montanari and Oh [37] and Chatterjee [22]. The survey [26] gives a global view of the field.

The dual problem of completion is sparsification, where given a matrix $P$, one seeks a procedure to keep only a handful of entries of $P$ without altering too much its properties [1, 27, 40, 49].

Those papers, although different in their methods, show that completing a matrix from the observation of $nd$ of its entries can only be done if the underlying matrix $P$ is not too complicated (i.e. low-rank and sufficiently incoherent), and in that case $P$ can efficiently be recovered only if $d$ is of order $\ln(n)$ — the so-called information-theoretic threshold for completion. In [37], there are results for $d$ fixed, but they are not sharp at all and do not allow any precise asymptotics on specific eigenvalues as we do. To our knowledge, the few works on completion from $d = O(n)$ entries (see for instance Gamarnik, Li and Zhang [32] and references therein) is focused on $\epsilon$-approximating the whole hidden matrix $P$, and never on exact estimation of a specific part of the matrix.

3.2 Random Matrices and Erdős-Rényi Graphs

When $P$ is the all-one matrix, the observed matrix $A$ is nothing but the renormalization of the adjacency matrix of a random Erdős-Rényi directed graph. Estimating the spectral properties of Erdős-Rényi undirected graphs is already quite difficult [38]. The complete description of the behavior of the greatest eigenvalues of Erdős-Rényi graphs have been totally explained, in the $d = o(n)$ sparse setting, only recently by different works: Benaych-Georges, Bordenave, Knowles [9, 10] and Alt, Ducatez and Knowles [2]. Recently, Tikhomirov and Youssef gave similar results for eigenvalues of Erdős-Rényi graphs with i.i.d. Gaussian weights on the edges [56]; here, the underlying matrix $P$ is thus drawn from GOE, and does not meet the usual assumptions of matrix completion. We finally mention a significant result on inhomogeneous Erdős-Rényi graphs by Chakrabarty, Chakraborty and Hazra [21] complementing [10].

In those works, it turns out that the behaviour of the (suitably normalized) high eigenvalues of Erdős-Rényi graphs is governed by the high degrees of the graph when $d < \ln(n)$, and stick to the edge $\pm 2$ of the limiting semi-circle law in then $d \to \infty$. The exact threshold for the disappearance of outliers happens at $d_* = \ln(4/e)^{-1} \ln(n)$ [2, 56]. Those results hold for undirected Erdős-Rényi graphs, and we are not aware of any similar results for directed Erdős-Rényi graphs, and even less in the really sparse regime where $d$ is fixed. Indeed, only the convergence of the global spectrum towards the circle law is now proven (when $d > \ln(n)^2$) by Basak and Rudelson [7]. Many questions and intuitions are given in the physicist survey [44]. Among them are listed (but not proved) our results on eigenvalues of Erdős-Rényi graphs, that will be stated in Sect. 4. Our results on eigenvectors completes the picture.
3.3 Phase Transitions

Our main result is a phase transition for the top eigenvalues of sparse non-Hermitian matrices: the whole bulk is confined in a circle of radius $O(1/\sqrt{d})$, and depending on the strength of the noise $d$, a few outliers appear and they are aligned with the corresponding eigenvalues of the original matrix $P$, and their eigenvectors have a nontrivial correlation with the original eigenvector.

This is of course similar to the celebrated BBP transition [6], and many similar transitions are already available in the literature of PCA or low-rank matrix estimation ([11, 41] and references therein). Apart from [45], which has a very different setting than ours, there are no results for phase transitions in low-rank non-symmetric matrix estimations, or in sparse settings and there are no results on fundamental limits in our setting.

A crucial question arising in this context is whether our methods are optimal in some sense, and if the thresholds we exhibit are sharp. In the litterature, there are not so many results on the impossibility of a statistical reconstruction task on sparse graphs, the most noticeable one being the impossibility of community detection on sparse SBM below the KS threshold [47]. Since our proof methods include the computation of martingale on random trees similar to the Kesten-Stigum martingale, an adaptation of [47] could be done at least in the setting where $P$ is a block-matrix. We leave this for future work.

On a side note, the comparison between non-backtracking and non-symmetric in Sect. 5 clearly indicate that the use of non-backtracking matrices detects more signal than the non-symmetric procedure, but nothing is known regarding the sharpness of this result. It would be interesting to study if ‘sparse completion problems’ like ours have an IT limit, and if there is a gap between the IT limit and a computational limit.

3.4 ‘Asymmetry helps’

One of the key features of this paper is that it deals with top eigenvalues of non-symmetric matrices. While the global behaviour of the spectrum of random matrices is now well understood (see the survey [13] on the circular law, or [33, 44] for physicist’s point of views), finer properties are less known.

Generally speaking, it is easier to deal with eigenvalues of Hermitian matrices, notably thanks to the variational characterizations of the eigenvalues. However, in many problems from applied mathematics, it turns out that the spectrum of Hermitian matrices can sometimes be less informative than the spectrum of other choices of non-Hermitian matrices. A striking instance of this fact was the so-called ‘spectral redemption conjecture’ in community detection [39] and [14], where the interesting properties were not captured by the spectrum of the adjacency matrix, but of a non-Hermitian matrix, the non-backtracking matrix.

In the setting of matrix perturbation, this insight was remarkably exposed in a recent and inspiring paper by Chen, Cheng and Fan [23]. Their setting is more or less the same as ours: an underlying Hermitian matrix $P$, which is asymmetrically perturbed into an observed non-Hermitian matrix $A = P + H$, the entries of $H$ being all i.i.d.
One might favor a singular value decomposition because of the conventional wisdom that SVD is more stable than eigendecomposition when it comes to non-Hermitian matrices; but this in fact not true, as shown in their Fig. 1, and indeed the eigenvalues are more accurate than the singular values; verbatim,

“When it comes to spectral estimation for low-rank matrices, arranging the observed matrix samples in an asymmetric manner and invoking eigen-decomposition properly (as opposed to SVD) could sometimes be quite beneficial.” [23, page 2]

This is the philosophy we would like to convey here; however, their result hold only on the not-so-sparse regime where \( d > \ln(n) \). We extend all their results to the fixed \( d \) regime, with an explicit threshold for the detection of \( P \) and exact asymptotics for perturbation of linear forms.

### 3.5 Eigenvalues of Perturbed Matrices

Many works on completion or sparsification rely on a perturbation analysis of the eigenvalues/singular values of perturbed matrices.

For example, one of the key points in many papers is that the sparsification procedure (from \( P \) to \( A \)) alters the spectral properties of \( P \), but not too much; indeed the top singular values or eigenvalues do not differ too much, hence keeping only the ‘greater’ items in the SVD or the eigendecomposition of \( A \) is sufficient to weakly recover \( P \); that was the idea of [22, 27, 37] (and many of their heirs). The proofs usually rely on estimates on eigenvalues/singular values of the random matrix \( A \), by combining concentration inequalities and eigenvalues inequalities (such as Weyl’s one). but no sharp asymptotics can be obtained with those methods, a limitation already visible in the seminal paper from Friedman, Kahn, Szemeredi [30] and Feige and Ofek ( [29]. This problem becomes unassailable when \( d \) is really smaller than \( \ln(n) \) or fixed, due to the fact that the underlying graphs are highly non-regular.

Our proof techniques globally rely on methods introduced by Massoulié and refined by Bordenave, Lelarge and Massoulié [14, 43]. This powerful and versatile trace method has now been used in various problems for estimating high eigenvalues of sparse random matrices, such as random regular graphs [12], biregular bipartite graphs [18], digraphs with fixed degree sequence [24], bistochastic sparse matrices [16], multigraph stochastic blockmodels [50]. However, our construction, and especially the pseudo-eigenvectors we chose, greatly simplifies the former analysis in [14, 43]. This considerable simplification has been very recently been applied to the non-backtracking spectrum of inhomogeneous graphs in Massoulié and Stephan in [54], it follow from our methods which was introduced in a preliminary version of this work contained in [25].

### 3.6 Eigenvectors of Perturbed Matrices

Eigenvector perturbation has also attracted a lot of attention, mainly around variants of the Davis-Kahan theorem [61]. As mentioned in [28], many algebraic bounds (such as Weyl’s inequality or the Davis-Kahan theorems) are tight in the worst case, but
wasteful in typical cases. Our proof method does not rely on those general bounds, and naturally integrates the perturbation of eigenvectors in combination with the now classical Neumann trick (see [23, 28]).

### 3.7 Extensions

The results in this paper extend naturally to the setting where the matrix $P$ is modeled as the Kronecker product of two matrices so that $P = P_A \otimes P_B$. In this setting, an application of the results of Van Loan and Pitsianis [59, Section 2] shows that there is a rearrangement operator $\mathcal{R}(P)$ which rearranges the elements of the matrix $P$ such that $\mathcal{R}(P)$ is rank one. In other words, $\mathcal{R}(P) = \text{vec}(P_A)\text{vec}(P_B)^*$ where vec(-) stacks the columns of the matrix on top of each other and creates a single column vector. We can use this to induce low rank matrices which fit into our framework. Kronecker product structured matrices are ubiquitous in many scientific applications (see [57, 58]) and spotting Kronecker structure in them and applying the re-arrangement trick can lead to improved matrix completion in Kronecker structured matrix completion problems.

Tensor completion (see [35]) is a natural extension of our ideas in matrix completion for the problem of completing higher order arrays with missing data. Low-rank higher-dimensional tensors can, via a rearrangement and flattening operation (see [34, Section 2]) be expressed as low-rank matrices (with not necessarily orthogonal components). This allows us to connect the low-rank incoherent tensor completion problem with our framework.

Finally, the idea that random asymmetrization bypasses the echo-chamber effect that cripples the SVD (or the symmetric eigen-decomposition) can be employed as a non-parametric estimation technique wherever the underlying singular vectors we are trying to estimate are delocalized but where the SVD (or a symmetric eigen-decomposition) returns localized estimates. This trick will work right out of the box, for example, in the estimation of low-rank matrices with incoherent singular vectors contaminated with heavy tailed noise or just-sparse-enough-but-too-large outliers. Trimming techniques which precisely tune the threshold parameters based on precise structural information might have lower estimator errors than a non-parametric technique such as ours but the random asymmetrization trick will be more robust to errors due to a mismatch in the structure model for the outliers or heavy tails. A hint that the random asymmetrization is a useful technique can spring from analytical or numerical simulation insights as in Fig. 6 whenever one gleans that the operator norm (or the largest singular value) might be asymptotically unbounded but that the spectral radius (or largest eigenvalue in magnitude) is not. Where else might this be trick be useful beyond our context? An important extension of our framework is in settings where the missing entries are not sampled uniformly. One important scenario, that lends significant structure in the pattern of the observed entries, corresponds to the setting where the entries are observed via (for example) a Poissonian process with an intensity that is proportional to a (assumed, known) function of the (magnitude of the) underlying matrix entries we are trying to estimate, as in for example [5].
We leave related explorations and excursions to follow-up work and interested readers.

## 4 Detailed Results: Square Matrices

In this section, we restrict ourselves to the case where \( P \) is a Hermitian square matrix of size \( n \times n \). We write its spectral decomposition as

\[
P = \sum_{k=1}^{n} \mu_k \varphi_k \varphi_k^*,
\]

(4.1)

the real numbers \( \mu_k \) are the eigenvalues of \( P \), and \( \varphi_1, \ldots, \varphi_n \) is an orthonormal basis of eigenvectors. The eigenvalues are ordered by decreasing modulus:

\[
|\mu_1| \geq \cdots \geq |\mu_n| \geq 0.
\]

As above, \( M \) is a matrix with i.i.d. Bernoulli entries with parameter \( d/n \), and the observed matrix is

\[
A = \left( \frac{n}{d} \right) P \odot M.
\]

Our goal is to describe the behaviour of the high eigenvalues of \( A \).

### 4.1 Main Result

Our result will hold uniformly over a wide class of matrices that match the usual hypothesis from the literature: low (stable) rank with incoherence conditions. The goal is not really to restrict the range of applications, but to track the dependence of the error terms with respect to the parameters at stake (such as stable rank, measure of incoherence or spectral separations). We list these definitions which are central to this paper, and then we explain them in the subsequent remark.

#### Definitions 4.1 (Complexity parameters of \( P \))

Let \( P = \sum \mu_i \varphi_i \varphi_i^* \) be a square Hermitian matrix. The amplitude, stable rank and incoherence describe the complexity of the matrix \( P \).

1. **Amplitude parameter \( L \):**

\[
L = n \max_{x,y} |P_{x,y}|.
\]

(4.2)

Equivalently, it is the scaled \( L^1 \) to \( L^\infty \) norm of \( P \).

2. **Stable numerical rank \( r \):**

\[
r = \frac{\|P\|_F^2}{\|P\|^2} = \sum_{k=1}^{n} \frac{\mu_k^2}{\mu_1^2}.
\]
3. **Incoherence parameter** \((b, r_0)\): any scalar and integer \(r_0 \leq n\) such that for every \(k\) in \([r_0]\) we have

\[
\max_{x \in [n]} |\varphi_k(x)| \leq \frac{b}{\sqrt{n}}. \tag{4.3}
\]

The incoherence parameter measures how delocalized are the \(r_0\) dominant eigenvectors: since they have unit norm, it is clear that \(b \geq 1\). The \(r_0\) in question will be defined just after.

**Definitions 4.2** (Detection parameters of \(P\) and \(d\)) Let \(P = \sum \mu_i \varphi_i \varphi_i^*\) be a square Hermitian matrix and \(d > 1\) be a real number. The detection threshold, rank and gap describe what parts of \(P\) can be detected and how easily.

1. **Variance matrix** \(Q\):

\[
Q_{x,y} = n|P_{x,y}|^2 \quad \rho = \|Q\|. \tag{4.4}
\]

2. **Detection threshold** \(\vartheta\): any number \(\vartheta\) such that

\[
\vartheta \geq \max\{\vartheta_1, \vartheta_2\}, \tag{4.5}
\]

where the ‘theta parameters’ are defined by

\[
\vartheta_2 = \sqrt{\frac{\rho}{d}} \quad \text{and} \quad \vartheta_1 = \frac{L}{d}.
\]

3. **Detection rank** \(r_0\): number of eigenvalues of \(P\) which have modulus strictly larger than \(\vartheta\), i.e.

\[
|\mu_1| \geq \cdots \geq |\mu_{r_0}| > \vartheta \geq |\mu_{r_0+1}| \geq \cdots \geq |\mu_n|. \tag{4.6}
\]

4. **Detection hardness or gap** \(\tau_0\):

\[
\frac{\vartheta}{|\mu_{r_0}|} = \tau_0 \in (0, 1). \tag{4.7}
\]

It is the gap between \(\vartheta\) and the smallest eigenvalue above \(\vartheta\).

We observe that our threshold \(\vartheta\) can vary above \(\max\{\vartheta_1, \vartheta_2\}\). This is to allow an optimal application of our main theorem below. It is also interesting to note that the usual algebraic rank of \(P\) is an upper bound on \(r\). We can now state our main theorem.

**Theorem 4.3** Let \(P\) and \(A\) be as above. We define \(D = \max(2d, 1.01)\) and

\[
\ell = \left[ (1/8) \log_D(n) \right]. \tag{4.8}
\]
There exists a universal constant $c \geq 1$ such that if the inequality

$$C_0 := cr_0^4 L^{22} \ln(n)^{16} \leq b^{-1} \tau_0^{-\ell},$$

(4.9)

holds true then with probability greater than $1 - c n^{-1/4}$, the following event occurs:

1) Eigenvalues. There exists an ordering of the largest $r_0$ eigenvalues in modulus $\lambda_1, \ldots, \lambda_{r_0}$ of $A$ such that for all $i \in [r_0]$,

$$|\lambda_i - \mu_i| \leq C_0 \left| \frac{\delta}{\mu_i} \right|^{\ell} |\mu_i|,$$

(4.10)

and all the other eigenvalues of $A$ have modulus smaller than $C_0^{1/\ell} \delta$.

2) Eigenvectors. We denote by $\psi_i$ and $\psi'_i$ two unit right and left eigenvectors of $\lambda_i$ with positive scalar product. The relative spectral gap ratio at $\mu_i$ is defined as

$$\tau_{i,\ell} = 1 - \min_{j \in [n] \setminus \{i\}} |1 - (\mu_j / \mu_i)^{\ell}|.$$  

(4.11)

Then, for every $i \in [r_0]$ and $j \in [r]$, one has

$$\left| \langle \psi_i, \varphi_j \rangle - \frac{\delta_{i,j}}{\sqrt{\gamma_i}} \right| \leq \frac{C_0 \tau_{i,\ell}^{\ell}}{1 - \tau_{i,\ell}}.$$  

(4.12)

where $\gamma_i \geq 1$ is the deterministic number only depending on $d$ and $P$ defined by

$$\gamma_i := \sum_{s=0}^{\ell} \frac{\langle 1, Q^s \varphi_i \otimes \varphi_i \rangle}{(\mu_i^2 d)^s}.$$  

(4.13)

The overlap between eigenvectors satisfy

$$\left| \langle \psi_i, \psi_j \rangle - \frac{|\Gamma_{i,j}|}{\sqrt{\gamma_i \gamma_j}} \right| \leq \frac{C_0 \tau_{i,\ell}^{\ell}}{\sqrt{(1 - \tau_{i,\ell})(1 - \tau_{i,\ell})}}.$$  

(4.14)

where $\Gamma_{i,j}$ are real numbers defined by

$$\Gamma_{i,j} := \sum_{s=0}^{\ell} \frac{\langle 1, Q^s \varphi_i \otimes \varphi_j \rangle}{(\sigma_i \sigma_j d)^s}.$$  

(4.15)

Finally, the same bound (4.12)-(4.14) also hold for the unit left eigenvectors $\psi'_i$ and

$$\left| \langle \psi_i, \psi'_i \rangle - \frac{\delta_{i,j}}{\gamma_i} \right| \leq \frac{C_0 \tau_{i,\ell}^{\ell}}{1 - \tau_{i,\ell}}.$$  

(4.16)
We have stated this theorem for any matrix $P$ with parameters $b, r, r_0, \tau_0, d, \vartheta$ without mentioning any dependence on $n$. In fact, all those parameters can indeed depend on $n$ since our result is non-asymptotic and quantitative. A numerical value for the universal constant $c \geq 1$ could be extracted from the proof, even if it would not be very informative: we have used various crude bounds to arrive at a tractable constant $C_0$ and a readable proof. There are however various ways to improve the value of $C_0$, notably by decreasing very substantially the factor $L^{22}$ or $\ln(n)^{16}$. We have postponed this technical discussion in the final Sect. 18.

In Theorem 4.3, the value of the threshold $\vartheta \geq \vartheta_2 \lor \vartheta_1$ is free, it determines the eigenvalues above the threshold and the gap $\tau_0$. Theorem 4.3 is not trivial if $C_0$ is smaller than $\tau_0^{-\ell}$. In the typical situation where the parameters $b, r, r_0$ are $O(\ln(n)^c)$, for some constant $c > 0$, it happens if $\log_d(1/\tau_0) \gg \ln(\ln(n))/\ln(n)$. We note also that in this paper we only focus in the regime where $d$ is small, typically for $d = O(\sqrt{\ln(n)})$, where usual spectral methods on the symmetric matrices are not working. We have thus made no effort in obtaining an interesting error bound when $d$ is larger.

The threshold $\vartheta_2 = \sqrt{\rho/d}$ is the analog of the Kesten-Stigum bound in community detection, see [46], it is related to an intrinsic property on the existence of an eigenwave in a Galton-Watson tree with Poisson offspring distribution with parameter $d$, see Sect. 14. In most applications and simulations, $\vartheta_2$ is bigger than $\vartheta_1 = L/d$. There is however a regime where $d$ is very small and as a consequence, the actual threshold is $\vartheta_1$. This is the same phenomenon as the one uncovered in [24] — see the definition of $\hat{\rho}$ in Theorem 1 of that paper, see also [15] for a similar phenomenon. In our setting, as it is defined, the threshold $\vartheta_1 = L/d$ is often pessimistic. We have used it mostly for the readability of the proofs. There are ways to decrease the value of $L$ for most choices of matrices $P$. We have again postponed this technical discussion in Sect. 18.
Remark 4.4 We note that it is immediate to check that as $d$ grows, $\gamma_i - 1 \sim C/d$ where $C$ depends on $P$.

Remark 4.5 The coefficients $\gamma_i$ have a simple expression if $P$ is such that $\sum_y Q_{xy} = n \sum_y P_{xy}^2$ does not depend on $x$. Indeed, in this case, the vector $\mathbf{1}$ is the top eigenvector of $Q$ and we have $\sum_x Q_{xy} = \rho/n$. In particular, $Q^s \mathbf{1} = \rho^s \mathbf{1}$. From (4.13), for $i \in [r_0]$, we get

$$\gamma_i = \sum_{s=0}^\ell \left( \frac{\vartheta_2}{\mu_i} \right)^{2s} = 1 - \left( \frac{\vartheta_2/\mu_i}{1 - \vartheta_2/\mu_i} \right)^{2(\ell+1)}.$$ (4.17)

We conclude this section with an easy corollary of Theorem 4.3. It asserts that, above the threshold $\vartheta$, the average of the left and right unit eigenvectors of $A$ is a good estimate of the corresponding eigenvector of $P$.

Corollary 4.6 With the notation of Theorem 4.3, for $i \in [r_0]$, let $\hat{\varphi}_i = (\psi_i + \psi_i') / |\psi_i + \psi_i'|$ where $| \cdot |$ is the Euclidean norm of a vector. On the event of Theorem 4.3, we have for all $j \in [r]$,

$$\left| \langle \hat{\varphi}_i, \varphi_j \rangle \right| - \sqrt{2} \delta_{i,j} \sqrt{\gamma_i + 1} \leq \frac{4C_0 \tau_0 \ell}{1 - \tau_{i,\ell}}.$$ (4.18)

The above result states that weak recovery is feasible even in the regime where $d$ is of order 1 provided that $\vartheta$ itself is of order 1. This is in very sharp contrast with what would happen if the revealed entries were symmetric. As known in the literature (a general survey is given in Sect. 3.2 at page 10), the top eigenvalues would then be aligned with the high-degree vertices, but also the top eigenvectors would be localized on those vertices, losing all the signal information.

4.2 Example: The Rank One Case and Erdos-Rényi Graphs

We illustrate Theorem 4.3 for rank one matrices: $P = \varphi \varphi^*$. With the above notation $r = 1$, $\mu_1 = 1$ and $\varphi_1 = \varphi$. In this case, from (4.17) it is easy to check that we have

$$\vartheta_2 = \sqrt{\frac{n |\varphi|_4^4}{d}}$$ and $$\gamma = \frac{1 - \vartheta_2^{2(\ell+1)}}{1 - \vartheta_2^2}.$$ (4.19)

where $\gamma = \gamma_1$ is defined by 4.13 and $|\varphi|_4^4 = \sum_x |\varphi(x)|^4$. For $d = O(\ln(n))$, Theorem 4.3 is an improvement of the results in [23].

This result on rank-one matrices can be applied to the adjacency matrix of a directed Erdős-Rényi graph. It corresponds to a matrix $P$ whose entries are $P_{x,y} = 1/n$ for all $x, y$. Then the matrix $A' = dA$ is the adjacency matrix of a random graph where each directed edge $(x, y)$ (including loops $(x, x)$) is present independently with probability $d/n$. In the asymptotic regime $n \to \infty$ and $d > 1$ is fixed, Theorem 4.3 implies that 1) $A'$ has one outlier eigenvalue close to $d$, all the other eigenvalues being smaller
Fig. 4 On top, the spectrum of a directed Erdős-Rényi graph with $d = 4$ an $n = 10000$ vertices. The outlier $\lambda_1 \approx 4$ is clearly visible. Above (right), a plot of the entries of the eigenvector of $A$ associated with $\lambda_1 \approx 4$. All entries are positive (Perron’s theorem) and are stacked close to the real eigenvector $\varphi$ which is in orange; their scalar product is close to $\sqrt{1 - 1/d}$. Above (left) is a histogram of the values of $\lambda_1$, over $N = 10000$ realizations of directed Erdős-Rényi graphs with parameters $n = 10000$ and $d = 4$.

than $\sqrt{d}$ and 2) the unit eigenvector $\psi$ associated with the outlier eigenvalue satisfies $|\langle \psi, 1/\sqrt{n} \rangle| = \sqrt{1 - 1/d} + o(1)$. These results are illustrated on Fig. 4, at page 16.

This is quite a striking contrast with the undirected sparse Erdős-Rényi graphs, where the high eigenvalues are aligned with the high-degree vertices, and the associated eigenvectors are localized on those vertices, see references below.

4.3 Example: Block-Matrix

In this section we work out a representative illustration of Theorem 4.3: a specific version of the SBM. We choose $P$ to be

$$P = \left( \begin{array}{cc} a & b \\ b & a \end{array} \right) \otimes E_{n/2} = \frac{2}{n} \begin{pmatrix} a & \ldots & a & b & \ldots & b \\ \vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\ a & \ldots & a & b & \ldots & b \\ b & \ldots & a & \ldots & b & \ldots \\ \vdots & \ddots & \vdots & \ddots & \ddots & \ddots \\ b & \ldots & b & a & \ldots & a \end{pmatrix}$$

(4.20)
with $E_k$ the matrix of size $k$ with 1 everywhere. Its spectral decomposition is given by $P = (a + b)1_11_1^* + (a - b)1_21_2^*$, where $1_1 = (1, \ldots, 1, 0, \ldots, 0)$ and $1_2 = 1 - 1_1$.

Such a $P$ obviously satisfies the required hypothesis for our analysis (low-rank, incoherence). The eigenvalues of $P$ are $a - b$, $a + b$ with multiplicity one, and 0 with multiplicity $n - 2$. Consequently, the matrix $Q = nP \odot P$ is equal to

$$
Q = \frac{4}{n} \begin{pmatrix}
a^2 & \ldots & a^2 & b^2 & \ldots & b^2 \\
\vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\
a^2 & \ldots & a^2 & b^2 & \ldots & b^2 \\
b^2 & \ldots & b^2 & a^2 & \ldots & a^2 \\
\vdots & \ddots & \vdots & \vdots & \ddots & \vdots \\
b^2 & \ldots & b^2 & a^2 & \ldots & a^2
\end{pmatrix},
$$

and we immediately infer $\rho = 2(a^2 + b^2)$. When $a = 4$ and $b = 1$, the nonzero eigenvalues of $P$ are thus 5 and 3 and the detection threshold is $\vartheta = \sqrt{34}/d$. When $d = 3$, this is approximately 3.317. Consequently, our theory says that $A$ will have one outlier eigenvalue close to 5, and all the other eigenvalues will be in the disk $D(0, 3.317)$, see Fig. 5 for an illustration.

5 Detailed Results: Non-Backtracking Matrix

In this section, we work out what happens if instead of using the adjacency matrix of the problem, we use the non-backtracking matrix.

5.1 Setting: Weighted Non-Backtracking Matrix

In the square symmetric case $P = P^*$, Theorem 4.3 and its Corollary 4.6 illustrate the striking accuracy of the spectrum of the non-symmetric matrix $A$ to estimate the symmetric matrix $P$. There is however some information which has been lost: the fact that $P = P^*$ has not been used in the definition of $A$ and the set of revealed entries could be almost doubled in principle. At some extra computational cost, a weighted variant of the non-backtracking matrix can cope with this issue.

Let us first define the probabilistic model. Let $\bar{M} \in M_n(\mathbb{R})$ be a random symmetric matrix where all entries above the diagonal are independent Bernoulli random variable with parameter $\bar{d}/n$: for all $x, y \in [n]$, $\bar{M}_{x,y} = \bar{M}_{y,x}$ and

$$
P(\bar{M}_{x,y} = 1) = 1 - P(\bar{M}_{x,y} = 0) = \frac{\bar{d}}{n}.
$$

Then, we define $E = \{(x, y) : \bar{M}_{x,y} = 1\}$, it is the set of revealed entries of $P$. With $\bar{d} = 2d - d/n$, this would correspond to the revealed entries of the matrix $(A + A^*)/2$ in our previous model up to a slight modification of the law of entries on the diagonal which is harmless in our setting.
We will need the following notation. A vector $\varphi \in \mathbb{R}^n$ can be lifted as two vectors, $\varphi^+, \varphi^-$ in $\mathbb{R}^E$ by setting $\varphi^-(x, y) = \frac{\varphi(x)}{\sqrt{\bar{d}}}$ and $\varphi^+(x, y) = \frac{\varphi(y)}{\sqrt{\bar{d}}}$.

The scaling is chosen so that these lifts are isometries from $\mathbb{R}^n$ to $L^2(\mathbb{R}^E, P)$:

$$E[|\varphi^\pm|^2] = |\varphi|^2,$$

where $|\cdot|$ is the Euclidean norm of a vector. The norm of $\varphi^\pm$ is tightly concentrated: if $|\varphi| = 1$ and $|\varphi|_\infty \leq b/\sqrt{n}$, then with probability at least $1 - 1/n$,

$$\left| |\varphi^+|^2 - 1 \right| \leq cb^2 d \ln(n)^{5/2} n^{-1/2},$$

for some universal constant $c > 0$ (it follows for example from the forthcoming Theorem 12.5).

The weighted non-backtracking matrix $B \in \mathcal{M}_E(\mathbb{R})$ is the non-symmetric matrix indexed by $E$ with entries, for $e = (x, y) \in E$ and $f = (a, b) \in E$ (those are directed edges):

$$B_{e,f} = \frac{n}{\bar{d}} 1_{a=y} 1_{x \neq b} P_{a,b}.$$

Exactly as for the matrix $A$, we can relate the top eigenvalues and eigenvectors of $B$ with those of $P$. The weighted non-backtracking matrix $B$ is defined on the directed edges of the graph induced by the non-missing entries of the matrix $P$ and hence, so are its eigenvectors. At each vertex, we sum the elements of an eigenvector of $B$ over all its incoming edges and then normalize the resulting vector to have unit norm. We refer to these vectors as the **weighting non-backtracking eigenvectors** of the matrix $P$.

### 5.2 Results

The non-backtracking matrix allows to reduce the detection threshold, the cost being that the size of the matrix $B$ is typically larger by a factor $\bar{d}$ (see Remark 5.6 below for possible ways to solve this issue).

A version of Theorem 4.3 also holds for the matrix $B$, but before stating it we need some definitions. The first ones are simply adapted from the square case. We emphasize the difference between the original definitions and the non-backtracking ones by overlining the corresponding quantities.

**Definitions 5.1 (Complexity parameters of $P$)** Let $P = \sum \mu_i \varphi_i \varphi_i^*$ be a square Hermitian matrix. The amplitude, stable rank and incoherence describe the complexity of the matrix $P$.

1. **Amplitude parameter $L$**:

$$L = n \max_{x,y} |P_{x,y}|. \quad (5.1)$$

---

---
Equivalently, it is the scaled $L^1$ to $L^\infty$ norm of $P$.

2. **NB-Stable rank $\tilde{r}$**: for technical reasons, we introduce a stronger notion of stable rank by setting

$$
\tilde{r} := \frac{\sum_{j=1}^{n} |\mu_j|}{|\mu_1|}.
$$

Note that $r \leq \tilde{r} \leq \text{rank}(P)$.

3. **Incoherence parameter** $(b, r_0)$: any scalar $b \geq 1$ and integer $r_0 \leq n$ such that for every $k$ in $[r_0]$ we have

$$
\max_{x \in [n]} |\varphi_k(x)| \leq \frac{b}{\sqrt{n}}.
$$

Again, the integer $r_0$ in the incoherence parameter is defined below to be equal to the detection rank.

**Definitions 5.2** (*NB-detection parameters of $P$ and $d$*) Let $P = \sum \mu_i \varphi_i \varphi_i^*$ be a square Hermitian matrix and $d > 1$ be a real number. The detection threshold, rank and gap describe what parts of $P$ can be detected and how easily.

1. **Variance matrix $Q$**:

$$
Q_{x,y} = n|P_{x,y}|^2, \quad \rho = \|Q\|.
$$

2. **NB-Detection threshold $\tilde{\vartheta}$**: any number $\tilde{\vartheta}$ such that

$$
\tilde{\vartheta} \geq \max(\tilde{\vartheta}_1, \tilde{\vartheta}_2),
$$

where

$$
\tilde{\vartheta}_2 = \sqrt{\frac{\rho}{d}} \quad \text{and} \quad \tilde{\vartheta}_1 = \frac{L}{d}.
$$

3. **Detection rank $r_0$**: number of eigenvalues of $P$ which have modulus strictly larger than $\vartheta$, i.e.

$$
|\mu_1| \geq \cdots \geq |\mu_{r_0}| > \tilde{\vartheta} \geq |\mu_{r_0+1}| \geq \cdots \geq |\mu_n|.
$$

4. **Detection hardness or gap $\tau_0$**:

$$
\frac{\tilde{\vartheta}}{|\mu_{r_0}|} = \tau_0 \in (0, 1).
$$

Before stating our main theorem for the non-backtracking matrix, we need to introduce a new parameter on the matrix $P$ which is really specific to the non-backtracking setting: this is due to the fact $B$ and $B^*$ are not equal in law.
Definition 5.3 If $r_0 \geq 1$, we define the matrix $C \in \mathcal{M}_{r_0}(\mathbb{R})$ by, for all $i, j \in [r_0]$,

$$C_{i,j} = \frac{\langle 1, Q\varphi_i \odot \varphi_j \rangle}{\mu_i \mu_j} = \sum_{x,y} \frac{Q_{x,y}}{\mu_i \mu_j} \varphi_i(x) \varphi_j(x).$$

Note that $C$ is scale invariant. The matrix $C$ is the Gram matrix of the vectors $(\varphi_i/\mu_i)_{i \in [r_0]}$ associated to the scalar product $\langle \psi, \phi \rangle_Q := \sum_x (\sum_y Q_{x,y}) \psi(x) \phi(x)$ on the vector space spanned by the coordinate vectors $(e_x)_{x \in V}$ where $V$ is the set of $x \in [n]$ such that $n \sum_y P^2_{x,y} = \sum_y Q_{x,y} > 0$. It is thus easy to check that $C$ is definite positive. We denote by $\sigma > 0$ the smallest eigenvalue of $C$. For example, if $\sum_y Q_{x,y}$ does not depend on $x$ then $\sigma \geq 1$. In general, we have $\sigma \geq \min_{x \in V} \sum_y Q_{x,y}/\mu_i^2$.

We are now ready to state a version of Theorem 4.3 for the non-backtracking matrix $B$.

Theorem 5.4 Let $P$ and $B$ be as above and assume $\bar{d} \geq 1$ and $\sigma \geq 0.01$. We define $D = \max(d, 1.01)$ and

$$\ell = \lceil(1/8) \log_D(n) \rceil. \quad (5.6)$$

There exists a universal constant $c \geq 1$ such that if the inequality

$$\bar{C}_0 = c d \bar{r}_0^4 L^{20} \ln(n)^{14} \leq b^{-1} \tau_0^{-\ell}, \quad (5.7)$$

holds true then with probability greater than $1 - cn^{-1/4}$, the following event occurs:

1) Eigenvalues. There exists an ordering of the largest $r_0$ eigenvalues in modulus $\lambda_1, \ldots, \lambda_{r_0}$ of $B$ such that for all $i \in [r_0]$,

$$|\lambda_i - \mu_i| \leq \bar{C}_0 \left| \frac{\delta}{\mu_i} \right|^{\ell} |\mu_i|, \quad (5.8)$$

and all the other eigenvalues of $B$ have modulus smaller than $\bar{C}_0^{1/\ell} \delta$.

2) Eigenvectors. We denote by $\psi_i$ and $\psi'_i$ two unit right and left eigenvectors of $\lambda_i$ with positive scalar product. For every $i \in [r_0]$ and $j \in [r]$, one has

$$\left| \langle \psi_i, \varphi_j^+ \rangle - \frac{\delta_{i,j}}{\sqrt{\gamma_i}} \right| \leq \bar{C}_0 \tau_0^{\ell} \frac{1}{1 - \tau_{i,\ell}} \quad \text{and} \quad \left| \langle \psi'_i, \varphi_j^+ \rangle - \frac{\delta_{i,j}}{\sqrt{\gamma_i}} \right| \leq \bar{C}_0 \tau_0^{\ell} \frac{1}{1 - \tau_{i,\ell}}. \quad (5.9)$$

where $\tau_{i,\ell}$ is defined in 4.11 and $\gamma_i$ is defined in 4.13 with $\bar{d}$ in place of $d$ and

$$\hat{\gamma}_i := \bar{d} \sum_{s=1}^{\ell+1} \frac{\langle I, Q^s \varphi_i \odot \varphi_j \rangle}{(\mu_i^2 \bar{d})^s}.$$
Finally, the same bound 4.12 also holds for the unit left eigenvector $\psi_i'$ and
\[
\left| \langle \psi_i, \psi_i' \rangle - \frac{\delta_i, j}{\sqrt{\gamma_i \gamma_j}} \right| \leq \frac{C_0 \tau_0^\ell}{1 - \tau_i, \ell}. \tag{5.10}
\]

The assumptions $\sigma \geq 0.01$ is only to guarantee a bound which is uniform in $\sigma \geq 0.01$. In general, it could easily be extracted from the proof an expression of $\bar{C}_0$ which depends on $\sigma$.

The coefficient $\hat{\gamma}_i$ has a simple expression if $\sum_y Q_{xy}$ is constant. Arguing as in Eq. 4.17, we have
\[
\hat{\gamma}_i = \tilde{d} \sum_{x=1}^{\ell+1} \left( \frac{\hat{\theta}_2}{\mu_i} \right)^{2s} = \frac{\rho}{\mu_i^2} \gamma_i = \frac{\rho}{\mu_i^2} \frac{1 - (\hat{\theta}_2/\mu_i)^2(\ell+1)}{1 - (\hat{\theta}_2/\mu_i)^2}. \tag{5.11}
\]

We will check that $\rho \geq \mu_i^2$ (in forthcoming 10.3). In particular, we always have in this case that $\hat{\gamma}_i \geq \gamma_i$. It follows from 5.9 that the right eigenvector $\psi_i$ is closer than the left eigenvector $\psi_i'$ to $\phi_i^+$. There is also an analog of Corollary 4.6 which allows to define a new sharp estimator of eigenvectors of $P$. To this end, we define the ‘left divergence’ of a vector $\psi \in \mathbb{R}^E$ as the vector $\tilde{\psi} \in \mathbb{R}^n$: for all $y \in [n],$
\[
\tilde{\psi}(y) = \frac{1}{\tilde{d}} \sum_{x:(x,y) \in E} \psi((x, y)).
\]
The ‘right divergence’ is defined as follows. Let $\deg(y)$ be the number of edges $E$ attached to $y$ (with loops, $\tilde{M}_{yy} = 1$, counting twice). If $\deg(y) \leq 1$, we set $\hat{\psi}(y) = 0$, otherwise, we set
\[
\hat{\psi}(y) = \frac{1}{d(\deg(y) - 1)} \sum_{x:(x,y) \in E} \psi((x, y)).
\]
With the notation of Theorem 5.4, we will check that, for symmetry reasons, we have for $i \in [r_0]$, the vectors $\hat{\psi}_i / |\hat{\psi}_i|$ and $\tilde{\psi}_i / |\tilde{\psi}_i|$ are very close to each other and well-defined.

**Corollary 5.5** With the notation of Theorem 5.4, for $i \in [r_0]$, let $\hat{\varphi}_i = \hat{\psi}_i / |\hat{\psi}_i|$ and $\tilde{\varphi}_i = \tilde{\psi}_i / |\tilde{\psi}_i|$. For some universal constant $c > 0$, with probability at least $1 - cn^{-1/4}$, we have for all $j \in [r],$
\[
\left| \langle \varphi_i, \varphi_j \rangle - \frac{\delta_i, j}{\sqrt{\gamma_i}} \right| \leq \frac{4C_0 \tau_0^\ell}{1 - \tau_i, \ell}.
\]
The same statement holds with $\tilde{\varphi}_i$ in place of $\hat{\varphi}_i$. 
Remark 5.6 The spectrum of $B$ is related through the so-called Ihara-Bass formulas to the spectrum of Hermitian matrices of dimension $n$, see [3, 60] for recent references. In the simplest case where the entries of $nP$ takes only two values, say 0 and 1, then the spectrum of $B$ can be obtained from the spectrum of a matrix in $\mathcal{M}_{2n}(\mathbb{R})$, see [4, Note 3.5]. These formulas have been used to design symmetric matrices in $\mathcal{M}_n(\mathbb{R})$ strongly connected to the spectrum of $B$, see notably [51, 52].

As an alternative, for an integer $\ell \geq 1$, we may consider the symmetric matrix $B_\ell \in \mathcal{M}_n(\mathbb{R})$ defined as

$$B_\ell = \nabla^* \Delta^{\ell-1} T \nabla,$$

where $\Delta \in \mathcal{M}_E(\mathbb{R})$ is the diagonal matrix defined for $e = (a, b) \in E$ by $(\Delta)_{e,e} = nP_{a,b}$, $T \in \mathcal{M}_E(\mathbb{R})$ is the involution matrix defined for $e = (a, b) \in E$ by $T\delta(a,b) = \delta(b,a)$ and the matrix $\nabla \in \mathcal{M}_{E,n}(\mathbb{R})$ defined for $e = (a, b) \in E$ and $x \in [n]$ by $\nabla_{e,x} = 1_{a=x}/\sqrt{d}$ (so that $\varphi^- = \nabla \varphi$). The entry $(B_\ell)_{x,y}$ is equal to the weighted sum of non-backtracking paths of length $\ell$ between $x$ and $y$ on the random graph whose adjacency matrix is $M$ and with edge weights $nP_{xy}/d$. Then with $\ell$ odd as in 5.6, it can easily be checked from the proofs that a version of Theorem 5.4 holds for $B_\ell$ if we replace the eigenvalues $\lambda_i$ of $B$ by $\text{sign}(\lambda_i,\ell) |\lambda_i,\ell|^{1/\ell}$ where the $\lambda_i,\ell$'s are the eigenvalues of $B_\ell$ (this result comes with better constants since we can rely on the spectral perturbation theory of symmetric matrices). In practice, this matrix $B_\ell$ is however less natural that the matrix $B$ since it has an extra parameter $\ell$ which is rather artificial.

5.3 Example 4.3 Continued

Let us go back to the example in Sect. 4.3 where $P$ was a block-matrix. We recall that the eigenvalues of $P$ are $a - b, a + b$ with multiplicity one, and 0 with multiplicity $n-2$. We also have $\rho = 2(a^2 + b^2)$. When $a = 4$ and $b = 1$, the nonzero eigenvalues of $P$ are thus 5 and 3 and the detection threshold with the non-symmetric masked matrix is $\vartheta = \sqrt{34/d}$. But now, the detection threshold with the non-backtracking matrix is $\tilde{\vartheta} = \sqrt{\rho/2d} = \sqrt{17/d}$. When $d = 3$, both thresholds will be

$$\vartheta = \sqrt{34/3} \approx 3.317 \quad \text{and} \quad \tilde{\vartheta} = \sqrt{17/3} \approx 2.23.$$

While the adjacency matrix will only have one outlier close to 5, the non-backtracking matrix will have two outliers, thus reflecting the whole structure of $P$ and capturing more information on $P$ — at a higher computational cost though, because the average size of $B$ is $2dn$. The phenomenon is illustrated at Fig. 5.

6 Detailed Results: Rectangular Matrices

We now go back to our original problem, where $P$ is a rectangular $m \times n$ matrix. Without loss of generality, we will assume $m \leq n$ and we introduce the parameter
Fig. 5 The underlying $1000 \times 1000$ symmetric matrix $P$ is as in Example 4.3 and has eigenvalues 0, 3 and 5. On the top panel, we see the eigenvalues of $A$ when the sparsity parameter is $d = 3$, with one outlier above the threshold $\widetilde{\vartheta} \approx 3.317$. Below is the spectrum of the non-backtracking matrix $B$ defined in this paragraph; there are two outliers close to 3 and 5, above the threshold $\widetilde{\vartheta} \approx 2.23$. Note that the second panel depicts around 6000 points (the average size of $B$ with $n = 1000$ and $d = 3$).
\[ \alpha \in (0, 1) \text{ defined as } \alpha = \frac{m}{n}. \tag{6.1} \]

Let \( M \) be a \( m \times n \) matrix whose entries are i.i.d. Bernoulli with parameter \( d/n \):

\[ P(M_{x,y} = 1) = 1 - P(M_{x,y} = 0) = \frac{d}{n}. \]

As before, the non-zero entries of \( M \) correspond to the entries of \( P \) that are observed.

For convenience, most proofs in this section are deferred to Sect. 17 at page 81.

### 6.1 Setting and Strategy: Reducing Non-Symmetric to Symmetric

We first describe a useful strategy to use our results for symmetric problems \( P \) and transfer them to the non-symmetric world.

We start by introducing an auxiliary \((n + m) \times (n + m)\) random matrix \( Z \) whose entries are Bernoulli random variables with the following distribution: for \( x, y \in [n + m] \),

\[
\begin{align*}
P(Z_{x,y} = 1, Z_{y,x} = 1) &= q & \text{ (} x \neq y \text{)} \tag{6.2} \\
P(Z_{x,y} = 0, Z_{y,x} = 1) &= P(Z_{x,y} = 1, Z_{y,x} = 0) = \frac{1 - q}{2} & \text{ (} x \neq y \text{)} \tag{6.3} \\
P(Z_{x,x} = 1) &= \frac{1 + q}{2} \tag{6.4}
\end{align*}
\]

where \( q \in (0, 1) \) is a parameter which must satisfy the following identity:

\[ q = \frac{d}{n} \left( \frac{1 + q}{2} \right)^2. \tag{6.5} \]

It is easy to check that such a \( q \) exists. It is given by \( 2d(1 - \sqrt{1 - d/n})/n - 1 \approx d/4n \). We finally define a \((n + m) \times (n + m)\) matrix with zero-one entries by

\[ \tilde{M} = Z \odot \begin{pmatrix} 0 & M^* \\ M & 0 \end{pmatrix}. \tag{6.6} \]

**Lemma 6.1** If \( q \) satisfies 6.5 then the entries of \( \tilde{M} \) are independent Bernoulli random variables with parameter

\[ \frac{2q}{1 + q} = (1 + o(1)) \frac{d}{2n}. \]

The proof is at Sect. 17.1. The mask \( \tilde{M} \) is thus an i.i.d. Bernoulli matrix with parameter \( 2q/(1 + q) = d/2n + o(1/n) =: \tilde{d}/\tilde{n} \) with \( \tilde{n} := n + m \). The parameter \( \tilde{d} \) is close to \((1 + \alpha)d/2\), hence we have \( \tilde{n}/\tilde{d} \approx 2n/d \). We can now apply our results from the first
section, especially Theorem 4.3, to our new estimator $\tilde{A}$ which we define now. First, we note $\tilde{P}$ the Hermitization of $P$:

$$\tilde{P} = \begin{pmatrix} 0 & P^* \\ P & 0 \end{pmatrix}.$$  \hspace{1cm} (6.7)

The link between $P$ and $\tilde{P}$ (especially between the spectral decomposition of $\tilde{P}$ and the SVD of $P$) is well-known in the literature; we recall it at Sect. 17.2. Our estimator is simply going to be

$$\tilde{A} := \left( \frac{2n}{d} \right) \tilde{M} \odot \tilde{P}.$$  \hspace{1cm} (6.8)

It is a block matrix with the following form:

$$\tilde{A} = \begin{pmatrix} 0 & A_1 \\ A_2^* & 0 \end{pmatrix}.$$  \hspace{1cm} (6.9)

where $A_1, A_2$ are $m \times n$ real matrices. Note that all the information contained in the original problem is kept intact: each revealed entry of $P$ is present at least once (maybe twice) in this new estimator $\tilde{A}$.

### 6.2 Results

Just as before, we first gather the main definitions involved in our result. We emphasize the differences with the quantities in the preceding sections with a tilde.

**Definitions 6.2 (Complexity parameters of $P$)** Let $P = \sum \sigma_i \zeta_i \xi_i^*$ be an $m \times n$ matrix and $\tilde{P}$ is Hermitization as in 6.7. The amplitude, stable rank and incoherence describe the complexity of the matrix $P$.

1. **Size**: $\tilde{n} = m + n$.
2. **Amplitude parameter** $\tilde{L}$:
   $$\tilde{L} = \tilde{n} \max_{x,y} |P_{x,y}| = (1 + \alpha)L.$$  \hspace{1cm} (6.10)
   Equivalently, it is the scaled $L^1$ to $L^\infty$ norm of $P$.
3. **Stable numerical rank** $r$:
   $$r = \frac{\|P\|^2_F}{\|P\|^2} = \frac{\sum_{k=1}^n \mu_k^2}{\mu_1^2}. $$
4. **Incoherence parameter** $(b, \tilde{r}_0)$: any scalar $b \geq 1$ and integer $\tilde{r}_0 \leq n$ such that for every $k$ in $[\tilde{r}_0]$ we have
   $$\max_{x \in [n]} |\xi_k(x)|, |\zeta_k(x)| \leq \frac{b}{\sqrt{n}}.$$  \hspace{1cm} (6.11)
The integer $\tilde{r}_0$ in the incoherence is defined below to be equal to the detection rank.

**Definitions 6.3 (Detection parameters of $P$ and $d$)** Let $\tilde{P}$ be as in 6.7 and $d > 1$ be a real number. The detection threshold, rank and gap describe what parts of $\tilde{P}$ (and $P$) can be detected and how easily.

1. **Variance matrix $\tilde{Q}$:**
   \[
   \tilde{Q}_{x,y} = \tilde{n}|\tilde{P}_{x,y}|^2 \quad \tilde{\rho} = \| \tilde{Q} \|.
   \] (6.12)

2. **Detection threshold $\tilde{\vartheta}$:** any number $\tilde{\vartheta}$ such that
   \[\tilde{\vartheta} \geq \max\{\tilde{\vartheta}_1, \tilde{\vartheta}_2\},\] (6.13)
   where the ‘theta parameters’ are defined by
   \[
   \tilde{\vartheta}_2 = \sqrt{\frac{\tilde{\rho}}{d}} \quad \text{and} \quad \tilde{\vartheta}_1 = \frac{L}{d}.
   \] (6.14)

3. **Detection rank $\tilde{r}_0$:** number of singular values of $P$ which are strictly larger than $\tilde{\vartheta}$, i.e.
   \[
   \sigma_1 \geq \cdots \geq \sigma_{\tilde{r}_0} > \tilde{\vartheta} \geq \sigma_{\tilde{r}_0+1} \geq \cdots \geq \sigma_n.
   \] (6.15)

4. **Detection hardness or gap $\tilde{\tau}_0$:**
   \[
   \frac{\tilde{\vartheta}}{\sigma_{\tilde{r}_0}} = \tilde{\tau}_0 \in (0, 1).
   \] (6.16)

Let us make a few remarks on these definitions.

- It is clear that if $Q$ is the non-square matrix defined by $Q_{x,y} = n|P_{x,y}|^2$, and $\rho = n\|Q\|$, thus corresponding to the base problem, then we have $\tilde{\rho} = (1 + \alpha)\rho$.

  The thresholds in 6.14 thus satisfy
  \[
  \tilde{\vartheta}_2 = \sqrt{\frac{\rho}{d}} \approx \sqrt{\frac{2\rho}{d}} \quad \text{and} \quad \tilde{\vartheta}_1 = \frac{L}{d} \approx \frac{2L}{d}.
  \]

- We know (see the link between $P$ and $\tilde{P}$ at Sect. 17.2, and more precisely 17.2) that for any $k \in [\tilde{r}_0], |\varphi_k|_{\infty} \leq b/\sqrt{2} \leq b$.

- The number of eigenvalues of $\tilde{P}$ with modulus greater than $\tilde{\vartheta}$ is $2\tilde{r}_0$. Each singular value $\sigma_i$ gives rise to two eigenvalues of $\tilde{P}$ at $+\sigma_i$ and $-\sigma_i$, as recalled in Sect. 17.2.

- It is easy to see that the stable rank $\tilde{r}$ of $\tilde{P}$ is equal to $2r$, where $r$ is the stable rank of $P$.

As in the square case, we need to define ‘theoretical correlations’ between eigenvectors. They now depend on more parameters than before.
Definition 6.4 (Theoretical covariances) For $i, j \in [\tilde{r}_0]$ and for signs $\circ, \square$, the real numbers $\Gamma_{i,j}^{\circ, \square}$ are defined by

$$\Gamma_{i,j}^{\circ, \square} := \sum_{s=0}^{\ell} \langle 1, \tilde{Q}_s^\circ \varphi_i^\square \otimes \varphi_j^\circ \rangle_{(\square \sigma_i)(\circ \sigma_j)d^s}$$

(6.17)

where $\ell$ is the integer defined in 6.18 thereafter.

We will also use $\gamma_i$ as a shorthand for $\Gamma_{i,i}^{+, +}$. It is easily checked to be bigger than 1; however, there is no particular reason for $\Gamma_{i,j}^{+, -}$ to be bigger than 1 or even positive in general, and indeed one can verify that $\lim_{d \to \infty} \Gamma_{i,j}^{+, -} = -\delta_{i,j}/2$.

We are now ready to state our main theorem for rectangular matrices: it directly follows from an application of Theorem 4.3 in our setting.

Theorem 6.5 Let $P$ and $\tilde{A}$ as defined earlier. We define $D = \max(2\tilde{d}, 1.01)$ and

$$\ell = \left\lfloor (1/8) \log_D (\tilde{n}) \right\rfloor.$$  

(6.18)

There exists a universal constant $c \geq 1$ such that if the inequality

$$C_0 := c \tilde{r}_0^4 b^4 L^{20} \ln(\tilde{n})^{16} \leq \tilde{\tau}_i^{-\ell},$$

(6.19)

holds true then with probability greater than $1 - cn^{-1/4}$, the following event occurs:

1) Eigenvalues. There exists an ordering $\lambda_1, \ldots, \lambda_{\tilde{r}_0}$ of the $\tilde{r}_0$ eigenvalues of $\tilde{A}$ with greater modulus and positive real part, such that for all $i \in [\tilde{r}_0]$,

$$|\lambda_i - \sigma_i| \leq C_0 \left| \tilde{\gamma}_i \right|^{\ell} |\sigma_i|,$$

(6.20)

and these $\lambda_i$ are real. All the other eigenvalues with positive real part have modulus smaller than $C_0^{1/\ell} \tilde{\gamma}$.

2) Eigenvectors. We denote by $\psi_i^{\pm}$ the unit right eigenvectors of $\pm \lambda_i$. The relative spectral gap ratio at $\sigma_i$ is defined as

$$\tilde{\tau}_{i,\ell} = 1 - \min_{j \in [n]\setminus\{i\}} |1 - (\sigma_j/\sigma_i)^\ell|.$$  

(6.21)

Then, for every $i \in [\tilde{r}_0]$ and $j \in [\tilde{r}_0]$, and for every signs $\circ, \square \in \{+, -\}$, one has

$$|\langle \psi_i^\circ, \varphi_j^\square \rangle| - \frac{\delta_{i,j}}{\sqrt{Y_i}} \leq \frac{C_0 \tilde{r}_0^\ell}{1 - \tilde{\tau}_{i,\ell}}$$

(6.22)

and

$$|\langle \psi_i^\circ, \psi_j^\square \rangle| - \left| \Gamma_{i,j}^{\circ, \square} \right| \leq \frac{C_0 \tilde{r}_0^\ell}{\sqrt{(1 - \tilde{\tau}_{i,\ell})(1 - \tilde{\tau}_{j,\ell})}}$$

(6.23)
Finally, if \( \psi_{i, \text{left}}^{\pm} \) denotes the unit left eigenvector associated with \( \pm\lambda_i \) with the convention 1.3, then

\[
\left| \langle \psi_{i, \text{left}}^{\Box}, \psi_{\circ,i, \text{left}}^{\circ} \rangle - \frac{\delta_{i,j}}{\sqrt{Y_i Y_j}} \right| \leq \frac{C_0 \bar{\tau}_0}{1 - \bar{\tau}_i, \ell}. 
\] (6.24)

The main difference between this theorem and the original theorem for symmetric matrices lies in the threshold \( \bar{\delta} \) in (6.13–6.14). With our method, the singular values of \( P \) are detected only above the new threshold

\[
\max \left\{ \frac{2L}{d}, \sqrt{\frac{2\rho}{d}} \right\}
\]

which is strictly bigger than \( \max\{L/d, \sqrt{\rho/d} \} \), the original threshold.

### 6.3 Smaller, Square Matrices

The matrix \( \tilde{A} \) is a square matrix with size \( m + n \), which is bigger than \( m \) by a factor \( 1 + \alpha^{-1} \). This can result in a higher computational cost, but we can do better and restrict ourselves to smaller matrices. Starting from the block decomposition (6.9), we can use the alternative matrices

\[
X := A_1 A_2^* \quad Y := A_2^* A_1
\] (6.25)

which are square matrices of respective sizes \( m \) and \( n \). The following example shows in details how to obtain them directly from the observations.

**Example 6.6** (From raw data to \( X \) and \( Y \)) Suppose that the matrix where we store the observed entries is

\[
\begin{pmatrix}
0 & 2 & 0 & 4 \\
1 & 0 & 0 & 4
\end{pmatrix}
\]

the zeros meaning that the corresponding entry is not observed. The steps to form the matrices \( X \) and \( Y \) are as follows: first, for each revealed entry, flip a coin as in Lemma 6.1 and put the entry right or left:

\[
\begin{pmatrix}
0 & 2 & 0 & 4 \\
1 & 0 & 0 & 4
\end{pmatrix} = \begin{pmatrix}
0 & 2 & 0 & 0 \\
0 & 0 & 0 & 4
\end{pmatrix} + \begin{pmatrix}
0 & 0 & 0 & 4 \\
1 & 0 & 0 & 0
\end{pmatrix}
\]

Second, normalize by \( 2n/d \) to get \( A_1 \) and \( A_2^* \):

\[
A_1 = \frac{2n}{d} \begin{pmatrix}
0 & 2 & 0 & 0 \\
0 & 0 & 0 & 4
\end{pmatrix} \quad A_2 = \frac{2n}{d} \begin{pmatrix}
0 & 0 & 0 & 4 \\
1 & 0 & 0 & 0
\end{pmatrix}
\]
Finally, multiply them to get the square matrices $X$ and $Y$: for $X$, which has size $m = 2$,

$$X = A_1 A_2^* = \left(\frac{2n}{d}\right)^2 \times \begin{pmatrix} 0 & 2 & 0 & 0 \\ 0 & 0 & 4 & 0 \end{pmatrix} \times \begin{pmatrix} 0 & 1 \\ 0 & 0 \\ 0 & 0 \\ 4 & 0 \end{pmatrix} = \left(\frac{2n}{d}\right)^2 \times \begin{pmatrix} 0 \\ 16 \end{pmatrix}$$

and for $Y$ which has size $n = 4$:

$$Y = A_2^* A_1 = \left(\frac{2n}{d}\right)^2 \times \begin{pmatrix} 0 & 1 \\ 0 & 0 \\ 0 & 0 \\ 4 & 0 \end{pmatrix} \times \begin{pmatrix} 0 & 2 & 0 & 0 \\ 0 & 0 & 4 & 0 \end{pmatrix} = \left(\frac{2n}{d}\right)^2 \times \begin{pmatrix} 0 & 0 & 0 & 4 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 8 & 0 & 0 \end{pmatrix}.$$

The elementary properties of those matrices are gathered in the following lemma, whose proof is a mere verification.

**Lemma 6.7 (Structure of $\tilde{A}$)** Let $\lambda$ be a nonzero eigenvalue of $\tilde{A}$ associated with a unit right-eigenvector $u v^\top$.

Then, $-\lambda$ is an eigenvalue of $\tilde{A}$, and a unit right-eigenvector is given by $(-u) v^\top$. (6.26)

If $X := A_1 A_2^*$ and $Y := A_2^* A_1$, then we have $\tilde{A}^2 = \text{diag}(X, Y)$. Moreover,

$$\text{Sp}(X) = \text{Sp}(Y) = \{\lambda^2 : \lambda \in \text{Sp}(\tilde{A})\}.$$ (6.27)

If a nonzero eigenvalue $\lambda$ has multiplicity $m_\lambda$ in $\tilde{A}$, then $\lambda^2$ has also multiplicity $m_\lambda$ in $X$ and in $Y$. If (6.26) was a unit right-eigenvector of $\lambda$, then $u \neq 0$ and $u/|u|$ is a unit right-eigenvector of $X$ associated with $\lambda^2$. Similarly, $v \neq 0$ and $v/|v|$ is a unit right-eigenvector of $Y$ associated with $\lambda^2$.

We can thus detect the squares of the singular values of $P$, and the singular values themselves, by only looking at $X$ or $Y$. Equivalently, we can estimate the singular vectors $\zeta, \xi$ by only looking at the eigenvectors of $X$ or $Y$. To do this, fix $i \in [\tilde{r}_0]$. We will note $\chi_i$ a unit right-eigenvector of $X$ associated with the eigenvalue $\lambda_i^2$, and $\chi_i'$ a left eigenvector (recall convention 1.3). Similarly, we will note $\pi, \pi_i'$ for the eigenvectors of $Y$. We will need a variation of the quantities $\gamma_i$. We define

$$\zeta_i \triangledown \xi_j := \begin{pmatrix} \zeta_i \odot \xi_j \\ 0 \end{pmatrix} \quad \xi_i \triangledown \xi_j := \begin{pmatrix} 0 \\ \xi_i \odot \xi_j \end{pmatrix}.$$ (6.28)
Equivalently, \( \zeta_i \triangle \zeta_j = \phi^+_i \otimes \phi^-_j - \phi^+_i \otimes \phi^-_j \) and \( \xi_i \triangledown \xi_j = \phi^+_i \otimes \phi^-_j + \phi^+_i \otimes \phi^-_j \). Then, with these notations we recall the quantities in 2.9 and 2.10:

\[
\Gamma^\triangle_{i,j} := \sum_{s=0}^\ell \frac{\langle 1, \tilde{Q}_i \triangle \zeta_j \rangle}{(\sigma_i^2 d)^s} \quad \Gamma^{\triangledown}_{i,j} := \sum_{s=0}^\ell \frac{\langle 1, \tilde{Q}_i \triangledown \xi_j \rangle}{(\sigma_i^2 d)^s}
\]

(6.29)

and we set \( \gamma_i^\triangle = \Gamma^\triangle_{i,i} \) and \( \gamma_i^{\triangledown} = \Gamma^{\triangledown}_{i,i} \). It is straightforward to check that \( \gamma_i^\triangle + \gamma_i^{\triangledown} = 2 \gamma_i \) and that

\[
\begin{align*}
\Gamma^\triangle_{i,j} &= \Gamma^{+,+}_{i,j} - \Gamma^{+,-}_{i,j} \\
\Gamma^{\triangledown}_{i,j} &= \Gamma^{+,+}_{i,j} + \Gamma^{+,-}_{i,j}.
\end{align*}
\]

(6.30)

**Theorem 6.8** Let, \( X, Y \) be the matrices defined in 6.25. We place ourselves under the event of Theorem 6.5. Then, there exists an ordering \( \nu_1, \ldots, \nu_{\tilde{r}_0} \) of the \( \tilde{r}_0 \) eigenvalues with greater modulus of \( X \), such that

\[
|\sqrt{\nu_i} - \sigma_i| \leq \frac{C_0 \tilde{\tau}^\ell_0}{1 - \tilde{\tau}_{i,\ell}}
\]

(6.31)

Let \( \chi_i \) and \( \chi'_i \) be two unit right and left eigenvectors associated with \( \nu_i \) with positive scalar product, then

\[
|\langle \chi_i, \zeta_j \rangle| - \left| \frac{\delta_{i,j}}{\sqrt{\gamma_i}} \right| \leq \frac{C_0 \tilde{\tau}^\ell_0}{1 - \tilde{\tau}_{i,\ell}}
\]

(6.32)

and

\[
|\langle \chi_i, \chi_j \rangle| - \frac{\Gamma^\triangle_{i,j}}{\sqrt{\gamma_i^\triangle \gamma_j^\triangle}} \leq \frac{C_0 \tilde{\tau}^\ell_0}{\sqrt{(1 - \tilde{\tau}_{i,\ell})(1 - \tilde{\tau}_{j,\ell})}}
\]

(6.33)

\[
|\langle \chi_i, \chi'_j \rangle| - \frac{\delta_{i,j}}{\gamma_i^\triangle} \leq \frac{C_0 \tilde{\tau}^\ell_0}{\sqrt{(1 - \tilde{\tau}_{i,\ell})(1 - \tilde{\tau}_{j,\ell})}}.
\]

(6.34)

A similar statement holds for \( Y \), its eigenvalues \( \nu_i \) and its eigenvectors \( \pi_i \), in particular:

\[
|\langle \pi_i, \pi_j \rangle| - \frac{\Gamma^{\triangledown}_{i,j}}{\sqrt{\gamma_i^{\triangledown} \gamma_j^{\triangledown}}} \leq \frac{C_0 \tilde{\tau}^\ell_0}{\sqrt{(1 - \tilde{\tau}_{i,\ell})(1 - \tilde{\tau}_{j,\ell})}}
\]

(6.35)

\[
|\langle \pi_i, \pi'_j \rangle| - \frac{\delta_{i,j}}{\gamma_i^{\triangledown}} \leq \frac{C_0 \tilde{\tau}^\ell_0}{\sqrt{(1 - \tilde{\tau}_{i,\ell})(1 - \tilde{\tau}_{j,\ell})}}.
\]

(6.36)

Finally, if the orientation of eigenvectors is chosen so that \( \langle \chi_i, \chi'_i \rangle \geq 0 \), then \( \langle \chi_i, \zeta_i \rangle \) and \( \langle \chi'_i, \zeta_i \rangle \) have the same sign.
Here we chose \( P \) as in Example 6.9 with singular values 5 and 3. The singular values of the masked matrix, \( A \), are depicted with the gray lines. They are uninformative, contrary to our matrices \( X \) and \( Y \) defined in (6.25). The spectra of \( X \) and \( Y \) are identical except for the multiplicity of the value 0, and is depicted in the picture. With \( d = 50 \) the two outliers can be seen close to \( \sigma_1^2 = 25 \) and \( \sigma_2^2 = 9 \).

The proof is in Sect. 17 at page 81 and the result is illustrated at Fig. 6.

The preceding theorem completely describes the behaviour of the most informative parts in the spectral decomposition of the ‘smaller matrices’ \( A_2^*A_1 \) and \( A_1^*A_2 \). They will also be used later in the design of ‘optimal’ estimators of the hidden matrix \( P \).

The theoretical covariances \( \Gamma_{i,j}^\Delta \) can be difficult to compute in general, however we will see in the following sections that

(i) when the rank of \( P \) is one, which is in itself an important example in applications, all the computations can explicitly be done (see Sect. 7 after);

(ii) When they cannot explicitly be computed, then can consistently be estimated by the mere results of Theorem 6.8. This will be done in the next section.

**Example 6.9** We took \( P \) to be a 1000 \( \times \) 5000 matrix with SVD \( P = 5\xi_1\xi_1 + 3\xi_2\xi_2 \), the singular vectors being taken uniformly at random over the unit sphere. With this matrix, we have \( \rho \approx 170 \) and the threshold is given by

\[
\tilde{\theta}_2 = \sqrt{\frac{2\rho}{d}} \approx \frac{18.47}{\sqrt{d}}.
\]

The singular value 3 will only be detected if \( 3 > \tilde{\theta}_2 \) or equivalently if \( d > (18.47/3)^2 \approx 37.91 \). With \( d = 50 \), the two outliers \( \nu_1, \nu_2 \) of \( X \) clearly appear close to the locations \( \sigma_1^2 = 5^2 = 25 \) and \( \sigma_2^2 = 3^2 = 9 \), same thing for \( Y \).
6.4 Statistical Estimation Aspects

We will now use the results of the preceding section for the statistical estimation of the singular vectors of \( P \). Fix some \( i \in [\tilde{r}_0] \). In the spectrum of \( X \), there is one eigenvalue \( \nu_i \) close to \( \sigma_i^2 \). It gives rise to two unit eigenvectors, \( \chi_i \) on the right and \( \chi'_i \) on the left, which contain information on \( \zeta_i \). Similarly, the unit eigenvectors \( \pi_i, \pi'_i \) of \( Y \) associated with \( \eta_i \) contain information on \( \xi_i \).

We present two estimators for \( \zeta_i \) and two for \( \xi_i \): using one eigenvector without any modification (‘simple’), or averaging the two left and right eigenvectors:

\[
\hat{\zeta}_i^\text{sim} = \chi_i \\
\hat{\zeta}_i^\text{avg} = \frac{\chi_i + \chi'_i}{|\chi_i + \chi'_i|}
\]

\[
\hat{\xi}_i^\text{sim} = \pi_i \\
\hat{\xi}_i^\text{avg} = \frac{\pi_i + \pi'_i}{|\pi_i + \pi'_i|}.
\]

In the end, the second estimator (averaged) is more powerful than the first one (simple). We mention the performances of both of them mainly for completeness.

The following theorem gives the full correlation between these estimators themselves, as well as their performance at estimating \( \zeta_i, \xi_j \). First, we define

\[
c_{1,i}^{\text{sim}} = \frac{1}{\sqrt{\gamma_i^{\Delta}}} \\
c_{1,i}^{\text{avg}} = \frac{2}{\sqrt{\gamma_i^{\Delta} + 1}} \\
c_{2,i}^{\text{sim}} = \frac{1}{\sqrt{\gamma_i^{\nabla}}} \\
c_{2,i}^{\text{avg}} = \frac{2}{\sqrt{\gamma_i^{\nabla} + 1}}.
\]

**Theorem 6.10 (Statistical estimators)** We place ourselves on the high-probability event of Theorem 6.5. Let \( i, j \in [\tilde{r}_0] \). Then, we have

\[
|\langle \hat{\zeta}_i^\text{sim}, \zeta_j \rangle| = (1 + o(1))c_{1,i}^{\text{sim}}, \\
|\langle \hat{\zeta}_i^\text{avg}, \zeta_j \rangle| = (1 + o(1))c_{1,i}^{\text{avg}} \\
|\langle \hat{\zeta}_i^\text{sim}, \xi_j \rangle| = (1 + o(1))c_{2,i}^{\text{sim}}, \\
|\langle \hat{\zeta}_i^\text{avg}, \xi_j \rangle| = (1 + o(1))c_{2,i}^{\text{avg}}
\]

Moreover, these estimators satisfy:

\[
\langle \hat{\zeta}_i^\text{sim}, \hat{\zeta}_j^\text{sim} \rangle = (1 + o(1))\frac{\Gamma_{i,j}^{\Delta}}{\sqrt{\gamma_i^{\Delta} \gamma_j^{\Delta}}}. \tag{6.36}
\]

\[
\langle \hat{\zeta}_i^\text{avg}, \hat{\zeta}_j^\text{avg} \rangle = (1 + o(1))\frac{\Gamma_{i,j}^{\nabla} + \delta_{i,j}}{\sqrt{(\gamma_i^{\nabla} + 1)(\gamma_j^{\nabla} + 1)}}. \tag{6.37}
\]

The proof is in Sect. 17.4 at page 83. One can observe that \( c_{k,i}^\# \) goes to 1 when \( d \to \infty \), indicating that in the high-degree regime where \( d \) is high, total reconstruction is nearly achieved. Moreover, it is easy to see, using Definition 6.17, that

\[
c_{1,i}^{\text{sim}} \leq c_{1,i}^{\text{avg}} \leq 1 \quad \text{and} \quad c_{2,i}^{\text{sim}} \leq c_{2,i}^{\text{avg}} \leq 1. \tag{6.38}
\]
The elementary identity $|\hat{\zeta}^# - \zeta|^2 = 2(1 - \langle \hat{\zeta}^#, \zeta \rangle) \approx 2(1 - c_{1,i}^#)$ shows that the closer $c_{1,i}^#$ is to 1, the better the estimator. The meaning of the inequalities in 6.38 is that avg estimator is better than the other, as it incorporates more spectral information.

### 6.5 Non-Backtracking Matrices

For better performances at a higher computational cost, it is naturally also possible to use the non-backtracking matrix in conjunction with the Hermitization $\tilde{P}$ of the matrix $P$ defined in (6.7). The mask matrix is then

$$\tilde{M} = \begin{pmatrix} 0 & M \\ M^* & 0 \end{pmatrix}.$$ 

We can thus define the weighted non-backtracking matrix $B$ associated to this mask matrix with weights $\tilde{P}$ as in Sect. 5. Applying directly Theorem 5.4, we then obtain a non-backtracking version of Theorem 6.5. This can be used to do statistical estimation of the singular vectors as in Sect. 6.4. To avoid too much repetitions, we leave the details of the statements to the reader since they follow from exactly the same considerations than above.

### 7 The Rank-one Case: Complete Theory

In this part, we develop in greater detail our theory when the underlying problem $P$ has rank 1. In this case, many computations can explicitly be done without too much difficulty and provide a better understanding of the different parameters at stake.

We begin with the simple case when $P$ is Hermitian, and then illustrate our statistical results when $P$ is rectangular with rank 1.

#### 7.1 Warm-up: Symmetric Problems

Here, the first basic model is the rank-one symmetric completion problem already explored in Sect. 4.2; the underlying matrix is $P = \varphi \varphi^*$ and the main parameters $\vartheta_2$ and $\gamma$ we computed in 4.19. We gather these results in the following proposition.

**Proposition 7.1** Suppose that $P = \varphi \varphi^*$. Then, $\rho = n|\varphi|^4$. The detection threshold is given by

$$\vartheta_2 = \sqrt{\frac{n|\varphi|^4}{d}}. \quad (7.1)$$

The parameter $\gamma_1$ is given by

$$\gamma_1 = \frac{1 - (\vartheta_2^2)^{\ell+1}}{1 - \vartheta_2^2}. \quad (7.2)$$
\[ = \frac{1 + o(1)}{1 - \frac{n|\varphi|^4}{d}}. \]  

(7.3)

But now, the eigenvector \( \varphi \) is going to be taken at random among various distributions, a common model in the literature. More precisely, we take a family \( (B(x))_{x \in [n]} \) of i.i.d. random variables, we set \( S = \sum |B(x)|^2 \) and we define

\[ \varphi(x) = \frac{B(x)}{S^{1/2}}. \]

From Theorem 4.3 and Eq. 4.19, the phase transition in \( d \) for weak recovery is given by

\[ n|\varphi|^4 = n \sum_{x \in [n]} |B(x)|^4 \left( \frac{\sum_{x \in [n]} |B(x)|^2}{n} \right)^2 \]

\[ = \left( \sum_{x \in [n]} |B(x)|^2 \right)^{-2} \left( \sum_{x \in [n]} |B(x)|^4 \right) \]

(7.4) \hspace{1cm} (7.5)

which is easily computed using the Law of Large Numbers: if \( B \) is a generic random variable with the same distribution as each \( B(x) \), and having a finite fourth moment, then almost surely one has

\[ n|\varphi|^4 = \frac{E[|B|^4]}{E[|B|^2]^2} (1 + o(1)) \sim \text{Kurt}_B \]

where \( \text{Kurt}_B \) is the non-centered kurtosis, i.e. the ratio of the fourth moment to the squared second moment; when the distribution is centered this is the classical kurtosis, defined as the ratio of the fourth centered moment to the fourth power of the standard deviation. Of course, the Cauchy-Schwarz inequality tells us that \( \text{Kurt} \) is always greater than 1 and this bound is attained for random variables with constant modulus — in particular, for the centered \( \text{Ber}(1/2) \) distribution. Table 1 collects some values of the kurtosis.

Remark 7.2 As mentioned in the introduction, the real threshold is \( \max \{ \vartheta_1, \vartheta_2 \} \), and in this setting it is equal to \( \vartheta_2 \) if and only if \( \sqrt{\rho/d} > L/d \), which reduces to \( n|\varphi|^4 > L^2/d \). But in our rank-one models, if the sampling distribution of entries of \( \varphi \) is unbounded, then \( L \) will grow to \( \infty \), even if very slowly: for standard normal random entries, we will have \( |\varphi|_\infty \approx \sqrt{\log n/n} \) so \( L \approx \log n \) and our theoretical threshold should (asymptotically) be \( L/d \). However, we can actually bypass this limitation, using tools introduced in [15] and further explained in the last section of this paper, Sect. 18. The key here is that \( L \) might be replaced by an essential supremum \( L' \), which accounts for the maximum of the entries of \( P \) after deleting a small subset of these entries. The simulations suggest that these refinements do indeed confirm that \( \vartheta_2 \) is the right threshold of interest.
Table 1  Some values of the kurtosis, appearing in our threshold as $n|\varphi|_4^4 \sim \text{Kurt}_B$

| distribution $B$ | Kurt$_B$, asymptotic value of $n|\varphi|_4^4$ |
|------------------|------------------------------------------|
| Bi-sided exponential (Laplace), $f(x) \propto e^{-|x|}$ | 6 |
| Hyperbolic secant | 5 |
| Standard normal | 3 |
| Uniform on $[0, 1]$ | $9/5 \approx 1.8$ |
| Ber($c$) | $1/c$ |
| Centered Ber(1/2) | 1 |
| Generalized normal: $f(x) \propto e^{-|x|^\beta}$ | $\frac{\Gamma(5/\beta)\Gamma(1/\beta)}{\Gamma(3/\beta)^2}$ |

7.2 Numerical Validation of Theoretical Results

Figure 7 shows the agreement between the theoretical predictions in Proposition 7.1 for the inner product between the left and right eigenvector and the ground truth vector and experiment for the setting where the rank one $P$ with $n = 8000$ has (unit-norm) eigenvectors vectors drawn from the uniform, normal and hyperbolic secant distributions.

The detection threshold, as in Proposition 7.1 is a function of the kurtosis and Fig. 7 confirms this prediction as well as that of the predicted inner products of the left/right eigenvectors with respect to the ground truth eigenvector and the prediction for the improved performance of the eigenvector estimate formed by averaging the left and right vectors.

The third column of Fig. 7 illustrates the improved accuracy of the estimated vectors in the $\ell_\infty$ error sense relative to the eigenvector obtained using the symmetric eigendecomposition. This goes beyond the statement of our results and an analysis of this improvement represents a natural follow-up of our line of work.

Note, too that the estimation performance gap between the asymmetric method and the symmetric method increases with the kurtosis of the eigenvector element. Intuitively this has to do with the localization of the eigenvectors of the symmetric matrix in the very sparse regime which is bypassed when we induce asymmetry.

7.3 Using Non-Backtracking Matrices

We illustrate in this section the behaviour of the non-backtracking statistics from Sect. 5, and especially Corollary 5.5. To do this, we recall that we first symmetrize the observation and then build the non-backtracking matrix $B$.

Working at the non-backtracking level. The left/right eigenvectors associated with the unique outlier of $B$ will be called $\psi, \psi'$, and lives in $\mathbb{C}^m$. The result in (5.9) says that if $\varphi^+$ is the lifting of $\varphi$, then

$$\langle \psi, \varphi^+ \rangle = \frac{1 + o(1)}{\sqrt{\gamma_i}},$$
Fig. 7  Numerical validation of the predictions in Proposition 7.1 for the setup described in Sect. 7.2 where \( n = 8000 \) and various values of \( d \)
just as in the preceding paragraph. The main difference now is that the left/right inner product is given thanks to (5.10) by

$$\langle \psi, \psi' \rangle = \frac{1 + o(1)}{\sqrt{\gamma_i \hat{\gamma}_i}}. \quad (7.6)$$

However, when the rank is 1, the quantity $\hat{\gamma}_i$ is found in (5.11) to be equal to $(\rho/\mu_1^2)\gamma_i = \rho \gamma_i$, which in our case is exactly

$$\hat{\gamma}_i = n|\varphi|_4^4 \gamma_i.$$

Consequently, the inner left/right product is given by

$$\langle \psi, \psi' \rangle = \frac{1 + o(1)}{\gamma_i \sqrt{n|\varphi|_4^4}} = \frac{1 - \frac{n|\varphi|_4^4}{d}}{\sqrt{n|\varphi|_4^4}} + o(1). \quad (7.7)$$

Asymptotically, the left and right non-backtracking eigenvectors are thus far from being aligned even when $d \to \infty$, since in this regime their angle converges towards $1/\sqrt{n|\varphi|_4^4}$ which is generally strictly smaller than 1 as soon as $\varphi$ is not the constant unit vector.

**Working with lowered eigenvectors** As in Corollary 5.5, we can also ‘lower’ the eigenvectors to the dimension $n$. To do this we simply follow one of the procedures described above Corollary 5.5; the inner product between the estimators $\hat{\varphi}$ and $\check{\varphi}$ and the real ground-truth eigenvector $\varphi$ is $1/\sqrt{\gamma_i}$ as above.

Figure 8 shows agreement between theory and experiment using the weighted non-backtracking matrix. The leftmost subplot in Fig. 8a confirms the accuracy of the inner product prediction in Theorem 5.4, and the equivalent performance of the left and right lowered vectors of the weighted non-backtracking matrix as predicted in Corollary 5.5. The rightmost subplot in Fig. 8a confirms the accuracy of the theoretical prediction for the inner product between the left and right (raised) eigenvectors of the weighted non-backtracking matrix given by Theorem 5.4 and (7.7) and that between the lowered left and right eigenvectors. The middle plot in Fig. 8a shows the improvement in eigenvector estimation due to the weighted non backtracking matrix relative to that obtained from the (symmetric) eigendecomposition and the average of the left and right vectors from the randomized asymmetric eigendecomposition. Figure 8b plots the same quantities as Fig. 8a, except over a single trial and with normally distributed eigenvectors – the plots confirm the accuracy of the asymptotic predictions and the concentration of measure implied in Theorem 5.4.

### 7.4 Rectangular Rank-One: Explicit Computations

Suppose now that that $P = \zeta \xi^*$ where $\zeta$, $\xi$ are unit vectors. The size of the matrix is $m \times n$ with $m/n = \alpha$. In this case, the whole problem relies on the computation of the quantities $\gamma_i^{\Delta \Delta}$. They might be difficult to compute in the general case, but here these
quantities can entirely be computed in terms of the 4-norm of $\zeta$, $\xi$, as in the symmetric case. The expressions are a little bit more intricate, but once $|\xi|_4$, $|\zeta|_4$ are known, the dependence in $d$ is simple. Note that $\gamma_i^{\Delta}$, $\gamma_i^{\nabla}$ both depend on $\ell$, see the definition in 6.17; however, in the computations (which are deferred to Sect. 17), this dependence can be neglected because it only gives rise to terms which are seemingly complicated, but who in the end behave like $\tilde{\theta}^\ell$ which goes to zero. This is why we encapsulated them in the $o(1)$ notation.

**Proposition 7.3** Suppose that $P = \zeta \xi^*$. The detection threshold is given by

$$
\tilde{\theta}_2 = \sqrt{\frac{2n|\zeta|^2_4|\xi|^2_4}{d}}.
$$

(7.8)
The parameters $\gamma^\triangle_i$, $\gamma^\triangledown_i$ are given by

$$
\gamma^\triangle_i = \frac{1 + o(1)}{2|\xi|^4} \left( \frac{|\zeta|^2 + |\xi|^2}{1 - \frac{2n|\zeta|^2|\xi|^2}{d}} + \frac{|\xi|^2 - |\zeta|^2}{1 + \frac{2n|\zeta|^2|\xi|^2}{d}} \right)
$$

(7.9)

and

$$
\gamma^\triangledown_i = \frac{1 + o(1)}{2|\xi|^4} \left( \frac{|\zeta|^2 + |\xi|^2}{1 - \frac{2n|\zeta|^2|\xi|^2}{d}} + \frac{|\xi|^2 - |\zeta|^2}{1 + \frac{2n|\zeta|^2|\xi|^2}{d}} \right).
$$

(7.10)

The proof of Proposition 7.3 is only a computation, although more tedious than when the underlying problem $P$ is Hermitian. We postponed it in Sect. 17, at page 83.

### 7.5 Rectangular Rank-One: Numerical Validation

Here, $P = \zeta\xi^*$, but we generate $\zeta$, $\xi$ using the same model as for the symmetric case: we put

$$
\zeta(x) = \frac{A(x)}{S_A^{1/2}} \quad \xi(x) = \frac{B(x)}{S_B^{1/2}}
$$

where $S_A = \sum |S(x)|^2$ and $S_B = \sum |B(y)|^2$, and $A(x)$ are i.i.d. samples from a common distribution $A$ and $B(x)$ are i.i.d. samples from another distribution. We suppose that both of them have finite fourth moments. With this model, the Law of Large Numbers entails

$$
S_A \sim nE[A^2] \quad \quad S_B \sim n\alpha E[B^2]
$$

$$
|\zeta|^4 \sim \frac{1}{n} \frac{E[A^4]}{E[A^2]^2} = \frac{\text{Kurt}_A}{n} \quad \quad |\xi|^4 \sim \frac{1}{\alpha n} \frac{E[B^4]}{E[B^2]^2} = \frac{\text{Kurt}_B}{\alpha n}.
$$

almost surely as $n \to \infty$.

In this case, Proposition 7.3 say that the detection threshold in $d$ is equal to

$$
2n|\zeta|^2|\xi|^2 = (1 + o(1))2\sqrt{\frac{\text{Kurt}_A\text{Kurt}_B}{\alpha}}.
$$

For example, if $A$, $B$ have the same kurtosis $k$, then the threshold for the birth of outliers close to 1 in the spectra of $X$ and $Y$ (defined in 6.25) is $2k/\sqrt{\alpha}$.

More precisely, with high probability, the following happens.

(i) If $d \leq 2n|\zeta|^2|\xi|^2$, then all the eigenvalues of $X$ and $Y$ have modulus smaller than $\sqrt{2n|\zeta|^2|\xi|^2/d} + o(1)$. 

\[\square\] Springer
(ii) If $d > 2n|\xi_4^2|\xi_4^2$, then all the eigenvalues of $X$ and $Y$ have modulus smaller than $\sqrt{2n|\xi_4^2|\xi_4^2/d + o(1)}$, except one eigenvalue $\nu$ of $X$ with $\nu = 1 + o(1)$ and one eigenvalue $\eta$ of $Y$ with $\eta = 1 + o(1)$.

We denote by $\chi, \chi'$ the unit right and left eigenvectors of $X$ associated with the outlier, when it exists. Similarly, we denote $\pi, \pi'$ the right and left eigenvectors of $Y$. Recall the convention 1.3 on the positivity of the scalar product of left and right eigenvectors.

**Estimators: definition and accuracy**

We will use the estimators defined in Sect. 6.4:

$$\hat{\zeta}_{\text{sim}} = \chi$$

$$\hat{\zeta}_{\text{avg}} = \frac{X + X'}{|X + X'|}$$

and similar estimators for $\xi$. We computed $\gamma_i^\triangle$ in Proposition 7.3: when $\text{Kurt}_A = \text{Kurt}_B = k$, a few manipulations show that

$$\gamma_i^\triangle \approx 1 + \frac{2k}{d} - \frac{4k^2}{d^2\alpha}$$

and

$$\gamma_i^\triangledown \approx 1 + \frac{2k}{d} - \frac{4k^2}{d^2\alpha}.$$  \hspace{1cm} (7.11)

Consequently, from Theorem 6.8 that the left/right inner product is

$$\langle \chi, \chi_i' \rangle = \frac{1 + o(1)}{\gamma_i^\triangle}$$

$$\langle \pi, \pi_i' \rangle = \frac{1 + o(1)}{\gamma_i^\triangledown}.$$  \hspace{1cm} (7.12)

The formulas for the correlations in Theorem 6.10 give $(\hat{\zeta}^\#, \zeta) = c_1^\#$, and we get

$$|\langle \hat{\zeta}_{\text{sim}}, \zeta \rangle| \approx \sqrt{\frac{1 - 4k^2/d^2\alpha}{1 + 2k/d}}$$

$$|\langle \hat{\zeta}_{\text{avg}}, \zeta \rangle| \approx \sqrt{\frac{1 - 4k^2/d^2\alpha}{1 + k/d - 2k^2/d^2\alpha}}.$$  \hspace{1cm} (7.13)

Finally, the MSE (in the sense of Proposition 2.2) in this rank-one context considerably simplifies. Indeed, we have $\sigma_1 = 1$ and above the threshold $\tilde{r}_0 = 1$, so that

$$\text{MSE}_{\#}^* = 1 - (c_{1,i}^\# c_{2,i}^\#)^2 + o(1)$$

and using the definitions of $c_{k,i}^\#$ from above Theorem 6.10 we find

$$\text{MSE}_{\text{sim}}^* = 1 - \frac{1}{\gamma_i^\triangle \gamma_i^\triangledown} + o(1)$$
\[ \text{MSE}_{\text{avg}}^* = 1 - \frac{4}{(1 + \gamma_i^{\Delta})(1 + \gamma_i^{\nabla})} + o(1). \]

**Illustrations** Figure 9 shows the agreement between theory and experiment for the rectangular setting with respect to the predicted inner product between the averaged left and right eigenvector of the \( X \) (resp. \( Y \)) matrix corresponding to the largest real eigenvalue and the left (resp. right) singular vector of the underlying matrix. The plot also confirms our prediction in Sect. 2.2 that the accuracy of the left (resp. right) singular vector estimated thus with respect to the ground truth vector can be determined from the inner product between the left and right eigenvectors of \( X \) (resp. \( Y \)). This underpins the statistically optimal (in the MSE sense) data-driven matrix completion algorithm.

We could lower the MSE of the recovered by using the weighted non-backtracking variant of the method – it is computationally too expensive for the \( m, n \) and \( d \) values considered here.

### 8 An Algebraic Perturbation Lemma

We present an eigenvalue-eigenvector perturbation theorem, which extends some of the results from [14, Section 4] by taking into account the lack of normality of the structures at stake. We formulate this tool in a separate section because it can be of independent interest.

Let us first give a simple description of the result: if \( u_i, v_j \) are vectors such that \( \langle u_i, v_j \rangle \approx \delta_{i,j} \), then every matrix close to \( S = \sum \theta_i u_i v_i^* \) has eigenvalues close to the \( \theta_i \), provided the \( u_i \)’s are sufficiently well-conditioned. Moreover, if the \( \theta_i \) are well-separated, the corresponding right-eigenvectors of \( A \) are close to the \( u_i \). Theorem 8.2 quantifies this for eigenvalues of generic matrices and Theorem 8.6 quantifies this for eigenvectors of matrix powers. The novelty here is that the vectors \( u_i \) need not form an orthonormal family for the result to hold, and the same for the \( v_j \)’s.

We first recollect the Bauer-Fike theorem:

**Theorem 8.1** (Bauer-Fike, [8], [55], chapter IV) Let \( S \) be a diagonalizable matrix, \( S = P \Sigma P^{-1} \) with \( \Sigma = \text{diag}(\theta_1, \ldots, \theta_n) \) and let \( A = S + E \) be a matrix. Then, all the eigenvalues of \( A \) lie inside the union of the balls \( B(\theta_j, \varepsilon) \) where \( \varepsilon = \|E\|\|P\|\|P^{-1}\| \). Moreover, if \( J \subset [n] \) is such that

\[
\bigcup_{j \in J} B(\theta_j, \varepsilon) \cap \left( \bigcup_{j \not\in J} B(\theta_j, \varepsilon) \right) = \emptyset,
\]

then the number of eigenvalues (with multiplicities) of \( A = S + E \) inside \( \bigcup_{j \in J} B(\theta_j, \varepsilon) \) is exactly \( |J| \). In particular, for each \( i \in [n] \), if \( m_i \) is the number of distinct eigenvalues of \( \Sigma \) which are in the connected component of \( \bigcup_{j \in J} B(\theta_j, \varepsilon) \) containing \( \theta_i \), there exists an ordering of the eigenvalues \( \lambda_1, \ldots, \lambda_n \) of \( A \) such that

\[
|\lambda_i - \theta_i| \leq (2m_i - 1)\varepsilon.
\]
The last statement is usually stated with $m_i = n$ (as in Theorem [55, Theorem 3.3], see Fig. 10 for an illustration (and first paragraph p170 in [55] for further explanation).

In the first section, we prove our general eigenvalue perturbation lemma. In the second section, we give a variant for powers of matrices, that incorporates a control over arguments of complex eigenvalues, and which also contains the eigenvector perturbation result.
The eigenvalue $\lambda_i$ need not be in $B(\theta_i, \varepsilon)$ because there might be some overlap with the closest balls, such as for the green or yellow ones in the drawing. However, if $m_i$ is the number of distinct balls in the connected component $B(\theta_i, \varepsilon)$, the eigenvalue $\lambda_i$ will always be within distance $(2m_i - 1)\varepsilon$ of some $\theta_{\pi(i)}$ and $\pi$ permutation in $\mathcal{S}_n$: for instance in the picture above one could take $\pi = (234)$.

8.1 Eigenvalue Perturbation

Let $u_1, \ldots, u_r, v_1, \ldots, v_r$ be two families of nonzero vectors in $\mathbb{R}^n$. Let us note $U = (u_1, \ldots, u_r)$ and $V = (v_1, \ldots, v_r)$; those are real matrices with $n$ lines and $r$ columns. Our ‘nearly diagonalizable’ matrix will be $S = U\Sigma V^*$ with $\Sigma = \text{diag}(\theta_1, \ldots, \theta_r)$, the $\theta_i$ being complex numbers. The center of our investigations will be some real square matrix $A \in \mathcal{M}_n(\mathbb{R})$, not necessarily diagonalizable, but close to $S$ in operator norm. We make the following assumptions.

1. There is some $\eta > 0$ such that

$$\|A - S\| \leq \eta. \quad (8.1)$$

2. The matrices $U$ and $V$ are well-conditioned, in the following sense:

- For some $N \geq 1$ we have $\|U\| \leq N$ and $\|V\| \leq N$.
- For some $h > 0$ we have

$$\lambda_{\text{min}}(V^*V) \geq h \quad \text{and} \quad \lambda_{\text{min}}(U^*U) \geq h. \quad (8.2)$$

where $\lambda_{\text{min}}$ is the smallest eigenvalue.
- The matrices $U$ and $V$ are nearly pseudo-inverses: there is a $\delta \geq 0$ such that

$$\|U^*V - I_r\| \leq \delta. \quad (8.3)$$
Theorem 8.2 Set $\theta_i = 0$ for all $i \in [n]\setminus[r]$ and
\[
\varepsilon = 12N^3\left(\eta + \frac{4\sqrt{r}\delta N^4 \max_i |\theta_i|}{h}\right).
\]

Under the preceding assumptions, we can apply the conclusion of Theorem 8.1 to $A$ and $\varepsilon$.

Proof By homogeneity, we may assume that $\max_i |\theta_i| = \|\Sigma\| \leq 1$. Note also that $\|S\| \leq \|U\| \|\Sigma\| \|\Sigma\| \|V\| \leq N^2$ and $\|A\| \leq \|A - S\| + \|S\| \leq \eta + N^2$. Hence the maximal distance between any pair of eigenvalues of $A$ and $S$ is at most $2(\eta + N^2)$. It follows that the statement is trivial if the following inequality does not hold:
\[
\delta \leq \frac{h}{4N^2 \sqrt{r}}.
\]

We will thus assume that the above inequality hold. We begin by defining a matrix $\tilde{U}$ close to $U$, which is really a pseudo-inverse of $V$; this will be achieved thereafter in (8.6). To do this, we define the vector spaces
\[
H_i = \text{vect}(v_j : j \neq i).
\]

Since $\lambda_{\min}(V^*V) \geq h > 0$, $V$ has full rank, hence $H_i$ has dimension $r - 1$. The orthogonal projection on $H_i$ is given by
\[
P_{H_i}(w) = V_i (V_i^* V_i)^{-1} V_i^* w
\]
where $V_i$ is $V$ whose $i$-th column $v_i$ has been deleted. Note that $V_i^* V_i$ is a principal submatrix of $V^* V$, hence it is nonsingular itself; moreover, its eigenvalues interlace those of $V^* V$ and in particular, its smallest eigenvalue is greater than $h$ through (8.2); when taking the inverse, we get $\|((V_i^* V_i)^{-1})\| \leq 1/h$.

We now consider the vectors defined by $\tilde{u}_i := u_i - P_{H_i}(u_i)$ and
\[
\tilde{u}_i := \frac{\tilde{u}_i}{\langle \tilde{u}_i, v_i \rangle} = \frac{u_i - P_{H_i}(u_i)}{\langle u_i - P_{H_i}(u_i), v_i \rangle}.
\]

We set $\tilde{U} = (\tilde{u}_1, \ldots, \tilde{u}_r)$; we want to prove that $\tilde{U}$ is close to $U$ and that $\tilde{U}^* V = I_r$. Let $e_j$ denotes the $j$-th element of the canonical basis of $\mathbb{R}^n$. By (8.3), we have $|V^* u_i - e_i| = |V^* U e_i - I_r e_i| \leq \|V^* U - I_r\| = \|U^* V - I_r\| \leq \delta$, thus we also have $|V_i^* u_i|^2 = \sum_{j \neq i} |\langle v_j, u_i \rangle|^2 \leq |V^* u_i - e_i|^2 \leq \delta^2$, and finally
\[
|u_i - \tilde{u}_i| = |P_{H_i}(u_i)| = |V_i (V_i^* V_i)^{-1} V_i^* u_i| \leq \|V\| \|(V_i^* V_i)^{-1}\| \delta \leq \frac{\delta N}{h}.
\]
Moreover, \( \langle u_i, v_i \rangle - 1 \) is the \( i \)-th diagonal entry of \( U^*V - I_r \) and thus its modulus is smaller than \( \|U^*V - I_r\| \), so we have \( |\langle u_i, v_i \rangle - 1| \leq \delta \), and

\[
|\langle \tilde{u}_i, v_i \rangle - 1| \leq |\langle u_i, v_i \rangle - 1| + |\langle u_i, v_i \rangle - \langle \tilde{u}_i, v_i \rangle|
\leq \delta + |u_i - \tilde{u}_i||v_i|
\leq \delta (1 + N^2 h^{-1})
\leq 2\delta N^2 h^{-1},
\]

where we have used that \( h \leq \|V\|^2 \). When \( 0 \leq t \leq 1/2 \), we have \( |(1+t)^{-1} - 1| \leq 3t/2 \), thus, since \( \delta \leq h/(4N^2) \), we have

\[
|1/\langle \tilde{u}_i, v_i \rangle - 1| \leq 3\delta N^2 h^{-1}.
\]

We now write

\[
|\tilde{u}_i - u_i| = \left| \frac{\tilde{u}_i}{\langle \tilde{u}_i, v_i \rangle} - u_i \right|
\leq |\tilde{u}_i - u_i| + \left| \frac{1}{\langle \tilde{u}_i, v_i \rangle} - 1 \right| |\tilde{u}_i|
\leq \delta Nh^{-1} + 3\delta h^{-1}N^2\|U\|
\leq \delta Nh^{-1}(1 + 3N^2).
\]

The last term is bounded by \( 4\delta N^3 / h \), so using the elementary inequalities \( \|M\| \leq \|M\|_F \), we get our estimation expressing how \( \tilde{U} \) and \( U \) are close:

\[
\|\tilde{U} - U\| \leq \frac{4\sqrt{r}\delta N^3}{h}.
\] (8.5)

Finally, from the definitions of \( H_i \) and \( \tilde{u}_i \), we have \( (\tilde{U}^*V)_{i,j} = \langle \tilde{u}_i, v_j \rangle = 0 \) if \( i \neq j \), and \( \langle \tilde{u}_i, v_i \rangle = 1 \), a crucial fact which can also be written as

\[
\tilde{U}^*V = V^*\tilde{U} = I_r.
\] (8.6)

Together, (8.5)–(8.6) achieve our preliminary and show that \( \tilde{U} \) is the suitable pseudo-inverse of \( V \) which is close to \( U \). We now study the matrix \( S \).

We set \( \Sigma = \text{diag}(\theta_i) \) and \( \tilde{S} = \tilde{U} \Sigma V^* \); we have

\[
\|S - \tilde{S}\| \leq \|V\|\|\Sigma\|\|\tilde{U} - U\|
\leq \frac{4\sqrt{r}\delta N^4}{h} := \eta',
\]

where we have used that \( \|\Sigma\| \leq 1 \). We now claim that \( \tilde{S} \) is diagonalizable with eigenvalues \( \theta_1, \ldots, \theta_r \). Indeed, \( \tilde{u}_i \) is an eigenvector with eigenvalue \( \theta_i \neq 0 \), and every

\( \odot \) Springer
basis of \( \text{im}(V)^\perp \) is a family of eigenvectors associated with the eigenvalue zero. We note \( \bar{S} = P^{-1} \Sigma' P \) with \( \Sigma' = \text{diag}(\theta_1, \ldots, \theta_r, 0, \ldots, 0) \) and \( P \) its diagonalization matrix.

The matrices \( A \) and \( \bar{S} \) are close:

\[
\| A - \bar{S} \| \leq \| S - \bar{S} \| + \| A - S \| \leq \eta' + \eta.
\]

We may thus apply Bauer-Fike Theorem 8.1 to \( A \) and \( \bar{S} \) with \( \epsilon = (\eta + \eta')\| P \|\| P^{-1} \| \). It thus remains to prove \( \| P \|\| P^{-1} \| \leq 12N^3 \). In the remaining of the proof. We compute \( P \) and \( \| P \|\| P^{-1} \| \).

Let \( K = \text{span}(v_1, \ldots, v_r)^\perp = \text{im}(V)^\perp = \ker(V^*) \); the dimension of \( K \) is \( n - r \). Let us choose any orthonormal basis \( (w_{r+1}, \ldots, w_n) \) of \( K \) and set up \( P = (\bar{U}, W) \) where \( W \) is the \( n \times (n - r) \) matrix whose columns are the \( w_k \)'s. Then, the family \( (\bar{u}_1, \ldots, \bar{u}_r, w_{r+1}, \ldots, w_n) \) is a diagonalization basis for the matrix \( S \); more precisely, we have \( S\bar{u}_i = \bar{U} \Sigma V^* \bar{u}_i = \theta_i \bar{u}_i \), and \( S w_j = 0 \). We now claim that the inverse of \( P \) is given by

\[
P^{-1} = \begin{pmatrix} V^* & 0 \\ -W^* \bar{U} V^* + W^* & \end{pmatrix}.
\]

We can directly check this using the relations (8.6), the orthonormality relation \( W^* W = I_{n-r} \) and \( V^* W = 0 \), which stems from the choice of \( W \) as a basis for \( \ker(V^*) \). Indeed,

\[
\begin{pmatrix} V^* \\ -W^* \bar{U} V^* + W^* \end{pmatrix} P = \begin{pmatrix} V^* \\ -W^* \bar{U} V^* + W^* \end{pmatrix} \begin{pmatrix} \bar{U} & W \end{pmatrix} = \begin{pmatrix} V^* \bar{U} \\ -W^* \bar{U} V^* + W^* \bar{U} + W^* \bar{U} \end{pmatrix} = \begin{pmatrix} I_r & 0 \\ 0 & I_{n-r} \end{pmatrix} = I_n.
\]

To compute the condition number of \( P \) we use the elementary Lemma 8.3, stated hereafter. Clearly, \( \| W \| = 1 \), hence by the lemma

\[
\| P \| \leq \sqrt{2} \max(1, \| \bar{U} \|).
\]

For \( P^{-1} \) we note that \( -W^* \bar{U} V^* + W^* = W^*(I_n - \bar{U} V^*) \), hence \( \| -W^* \bar{U} V^* + W^* \| \leq \| W \|\| I_n - \bar{U} V^* \| \), and

\[
\| P^{-1} \| \leq \sqrt{2}(1 + \| \bar{U} \|\| V \|).
\]
We thus get
\[ \| P \| \| P^{-1} \| \leq 2 \max(1, \| \tilde{U} \|) (1 + \| \tilde{U} \| \| V \|) . \]

From (8.5), \( \| \tilde{U} \| \leq \| U \| + \| \tilde{U} - U \| \leq N + 4 \sqrt{r} \delta N^3 / h \leq 2N \) from our assumption on \( \delta \). Finally, we get
\[ \| P \| \| P^{-1} \| \leq 4N(1 + 2N^2) \leq 12N^3 . \]
It concludes the proof. \( \Box \)

As promised, here is a simple lemma used in the preceding proof.

**Lemma 8.3** Let \( M_1 \in M_{n,r}(\mathbb{R}) \) and \( M_2 \in M_{n,n-r}(\mathbb{R}) \) be two matrices; we set \( M = (M_1, M_2) \in M_{n,n}(\mathbb{R}) \). Then
\[ \| M \| \leq \sqrt{2} \max\{\| M_1 \|, \| M_2 \|\} . \]

**Proof** For any \( x \in \mathbb{R}^r, y \in \mathbb{R}^{n-r} \) and \( z = (x^*, y^*)^* \) we have
\[ |Mz| = |M_1x + M_2y| \leq \| M_1 \| |x| + \| M_2 \| |y| \]
\[ \leq \max\{\| M_1 \|, \| M_2 \|\} \sqrt{2} \sqrt{|x|^2 + |y|^2} \]
\[ = \max\{\| M_1 \|, \| M_2 \|\} \sqrt{2} |z| \]
which is valid for any \( z \in \mathbb{R}^n \). \( \Box \)

### 8.2 Eigenvector Perturbation

Perturbation theory for invariant subspaces is a delicate matter. For simplicity, we will restrict ourselves here to the case of invariant subspaces of dimension 1, that is eigenvectors of simple eigenvalues.

We consider the matrix \( A = S + E \) with \( S = U \Sigma V^* \) as in Theorem 8.2. If \( \Sigma = \text{diag}(\theta_1, \ldots, \theta_r) \), we define the spectral ratio gap at \( \theta_i \) as
\[ g_i = \min_{1 \leq j \neq i \leq n} |1 - \theta_j / \theta_i| . \]
(8.8)
where we have set \( \theta_i = 0 \) for all \( i \in [n] \setminus [r] \).

**Theorem 8.4** Under the assumptions of Theorem 8.2, we consider the ordering of the eigenvalues of \( A, \lambda_1, \ldots, \lambda_n \) defined in Theorem 8.2. Set
\[ \varepsilon = 12N^3 \left( \eta + \frac{4 \sqrt{r} \delta N^4 \max_i |\theta_i|}{h} \right) . \]
Then for any $i \in [r]$, if $g_i \theta_i > 2\varepsilon$, $\lambda_i$ is a simple eigenvalue and any unit right-eigenvector of $A$ associated with $\lambda_i$, denoted $\psi_i$, satisfies

$$
\left| \frac{\langle \psi_i, \frac{u_i}{|u_i|} \rangle}{|u_i|} \right|^2 \geq 1 - \frac{12 g_i^2 h^2}{|\min_{j \in [r]} |\theta_j|} + \frac{N^2 \eta}{|\theta_i| - 2\varepsilon},
$$

(8.9)

and the same result holds for a unit left-eigenvector and $v_i/|v_i|$.

The proof is inspired by the so-called Neumann trick. We start with an elementary statement.

**Lemma 8.5** Let $D = \text{diag}(d_1, \ldots, d_r)$ with $d_i$ complex numbers such that $\min_{j \neq 1} |d_1 - d_j| = g > 0$. If $w = (w_1, \ldots, w_r) \in \mathbb{C}^r$ is a unit vector then $|w|_1^2 \geq 1 - |d_1 w - Dw|^2 / g^2$.

**Proof** We have $1 - |w|_1^2 = |w_2|^2 + \cdots + |w_r|^2 \leq |(d_1 I_r - D) w|^2 / g^2$. □

**Proof of Theorem 8.4** We fix some $i \in [r]$, and we note $\lambda = \lambda_i$, $\theta = \theta_i$, and $g = g_i$. We set $E = A - S$, so that, by assumptions, $\|E\| < \eta$. By scale invariance, we may assume without loss of generality that $(\min_i |\theta_i|)^{-1} = \|\Sigma^{-1}\| = 1$. Since $g \theta > 2\varepsilon$, from Theorem 8.2, it implies that $|\lambda - \theta| < \varepsilon$ (the ball of radius $B(\theta, \varepsilon)$ does not intersect any of the balls of radius $B(\theta, \varepsilon)$ with $j \neq i$). In particular $|\lambda| > |\theta| - \varepsilon \geq \gamma - \varepsilon > \varepsilon > \eta$ and hence $\lambda$ is not an eigenvalue of $E$ and $\det(E - \lambda I_n) \neq 0$. Then, using Sylvester’s identity,

$$
0 = \det(A - \lambda I_n) = \det(U \Sigma V^* + E - \lambda I_n)
= \det(E - \lambda I_n) \det(I_n + U \Sigma V^*(E - \lambda I_n)^{-1})
= \det(E - \lambda I_n) \det(I_r + \Sigma V^*(E - \lambda I_n)^{-1} U)
$$

which implies that $\det(I_r + \Sigma V^*(E - \lambda I_n)^{-1} U) = 0$. Consequently, there is a vector $w \in \mathbb{C}^r$ with $|w| = 1$ such that

$$
0 = w + \Sigma V^*(E - \lambda I_n)^{-1} U w.
$$

(8.10)

We set $\phi = \lambda(E - \lambda)^{-1} U w$. This vector cannot be zero (because $w \neq 0$ and $\Sigma V^*$ is nonsingular), and it is actually an eigenvector of $A$ associated with the eigenvalue $\lambda$:

$$
(A - \lambda I_n) \phi = \lambda(U \Sigma V^* + E - \lambda I_n)(E - \lambda I_n)^{-1} U w
= \lambda(U \Sigma V^*(E - \lambda I_n)^{-1} U + U)w
= \lambda U(\Sigma V^*(E - \lambda I_n)^{-1} U + I_r)w = 0.
$$

Our goal is now to prove that $w$ is close to the vector $e_i$ of the canonical basis. We set $C = \lambda V^*(E - \lambda I_n)^{-1} U + I_r$. In (8.10), we see that $0 = \lambda w - \Sigma w + \Sigma C w$, hence $|\lambda \Sigma^{-1} w - w| \leq \|C\|$ and we want to bound $C$. This is done as follows:
From the Neumann series expansion,

$$
\| (E - \lambda)^{-1} - (-\lambda^{-1})I_n \| \leq \frac{1}{|\lambda|} \frac{\|E\|}{1 - \|E\|/|\lambda|}.
$$

We thus obtain the bound:

$$
\| C \| \leq \delta + \frac{N^2 \eta}{|\lambda| - \eta} =: \delta'.
$$

and we obtain $|\lambda \Sigma^{-1} w - w| < \delta'$. We thus have $|\theta \Sigma^{-1} w - w| < |\theta - \lambda| \Sigma^{-1} | + |\lambda \Sigma^{-1} w - w| < \varepsilon + \delta'$. We apply Lemma 8.5 to $D = \theta \Sigma^{-1}$ and we get that $|w_i|^2 \geq 1 - (\varepsilon + \delta')^2 / g^2$. We may assume without loss of generality that $w_i$ is a non-negative real number. We deduce that $|w - e_i|^2 = 2 - 2\|\theta(w_i)\| \leq 2(\varepsilon + \delta')^2 / g^2$. Then we find that

$$
|\phi - u_i| \leq \|\lambda(E - \lambda)^{-1} U - I_r\| + |w - e_i|
$$

$$
\leq \frac{\|U\| \|E\|}{|\lambda| - \|E\|} + \frac{\sqrt{2}(\varepsilon + \delta')}{g} =: \varepsilon'.
$$

But keep in mind that $\phi$ might not be normalized. A unit-norm right eigenvector of $A$ associated with $\lambda$ is $\phi/|\phi|$. However $|\phi| - |u_i| \leq |\phi - u_i|$, so that $|\phi/|\phi| - u_i/|u_i|| \leq 2|u_i - \phi/|u_i||$, and finally, using $|u_i| > h$, we get

$$
|\phi/|\phi| - u_i/|u_i|| \leq 2\varepsilon'/h.
$$

Simplifying the error term leads to the bound (8.14) (observe that $g \leq 1$ and any unit eigenvector $\psi$ associated with $\lambda$ satisfies $|(\psi, u_i/|u_i||)^2 = |(\phi/|\phi|, u_i/|u_i||)^2 = 1 - |\phi/|\phi| - u_i/|u_i|||^2/2$).  

\[\Box\]

### 8.3 Powers of Matrices

We now use the preceding perturbation theorems, but for powers of matrices. This will be necessary since the perturbations that we will encounter are not small in norm unless we raise them at a high power. We need some variations on the hypothesis which are tuned for our needs. We emphasize the fact that they are certainly not optimal, but rather suited to our subsequent needs. Let $U, U', V, V'$ be four $n \times r$ matrices and $\ell, \ell'$ be two integers. We set $S = U \Sigma^\ell V^*$ and $S' = U' \Sigma^{\ell'} (V')^*$, where $\Sigma = \text{diag}(\theta_1, \ldots, \theta_r)$ and the $\theta_i$’s are real numbers (unlike the previous paragraphs where they were complex).

Let $A \in \mathcal{M}_n(\mathbb{C})$. We make the following assumptions.

\(\square\) Springer
1. The integers $\ell$ and $\ell'$ are mutually prime.

2. There are numbers $\eta, \eta' > 0$ such that
   \[ \| A^\ell - S \| \leq \eta \quad \text{and} \quad \| A^{\ell'} - S' \| \leq \eta' \]  
   \hspace{1in} (8.11)

3. The matrices $U, U', V, V'$ are well-conditioned:
   - The operator norms of $U$ and $V$ are smaller than $N$.
   - The smallest eigenvalue of $U^* U, V^* V$ are greater than $h > 0$.
   - The operator norm of $U^* V - I_r$ is smaller that $\delta$.
   - The above properties hold for $U'$ and $V'$ for some constants $N', h', \delta'$.

The following perturbation theorem holds.

**Theorem 8.6** Under the above assumptions, we set

\[ \varepsilon = 12N^3 \left( \eta + \frac{4\sqrt{\delta}N^4 \max_{i} |\theta_i|^\ell}{h} \right) \]

and $\varepsilon'$ is defined similarly with the parameters $(\ell', \eta', N', h', \delta')$. Assume that for all $i \in [r]$,

\[ |\theta_i|^\ell > 2\ell' r \varepsilon \quad \text{and} \quad |\theta_i|^\ell' > 2\ell' r \varepsilon'. \]  
   \hspace{1in} (8.12)

Then, there is an ordering of the $r$ largest eigenvalues of $A$ in modulus, say $\lambda_1, \ldots, \lambda_r$ such that

\[ |\lambda_i - \theta_i| < \frac{4r \varepsilon}{\ell |\theta_i|^\ell - 1}. \]  
   \hspace{1in} (8.13)

All the other $n - r$ eigenvalues of $A$ have modulus smaller than $\varepsilon^{1/\ell}$.

Moreover, for any $i \in [r]$, if $g_i := \min_{j \neq i} |1 - \theta_j^\ell / \theta_i^\ell| > 2\varepsilon \theta_i^{-\ell}$, then $\lambda_i$ is a simple eigenvalue and any unit right-eigenvector of $A$ associated with $\lambda_i$, denoted $\psi_i$, satisfies

\[ \left| \langle \psi_i, \frac{u_i}{|u_i|} \rangle \right|^2 \geq 1 - \frac{12}{g_i^2 h^2} \left( \frac{\varepsilon}{\min_{j \in [r]} |\theta_j|^\ell} + \frac{N^2 \eta}{|\theta_i|^\ell - 2\varepsilon} \right)^2 \]  
   \hspace{1in} (8.14)

and the same result holds for a unit left-eigenvector and $v_i / |v_i|$.

**Proof** We start with the statement on eigenvalues. We apply Theorem 8.2 to $A^\ell$ and $S$ and to $A^{\ell'}$ and $S'$. There are two permutations $\pi, \pi' \in \mathfrak{S}_r$ such that $|\lambda_i^\ell - \theta_{\pi(i)}^\ell| \leq \varepsilon_0 = (2r - 1)\varepsilon$ and $|\lambda_i^{\ell'} - \theta_{\pi'(i)}^{\ell'}| \leq \varepsilon'_0 = (2r - 1)\varepsilon'$. Indeed, by assumption (8.12), the ball of radius $\varepsilon$ centered at 0 does not intersect any of the balls centered at $\theta_i, i \in [r]$.

We may assume that $\pi$ is the identity.
We fix some \( i \in [r] \). Our goal is to show that indeed, \( \lambda_i \) is close to \( \theta_i \). We first control the argument of \( \lambda_i \) by using that two mutually prime powers of \( \lambda_i \) are close to the real axis (since the \( \theta_j \)'s are real). More precisely, we set \( x = \theta_i \) and \( y = \theta_{\pi'(i)} \), so that

\[
|\lambda_i / x^\ell - 1| \leq u := \varepsilon_0 / |x|^\ell
\]

and

\[
|\lambda_i' / y^{\ell'} - 1| \leq u' := \varepsilon'_0 / |y|^{\ell'}.
\]

The polar decomposition of \( \lambda_i \) is written \( \lambda_i = |\lambda_i| e^{i\omega} \) with \( \omega \) real. The argument of \( (\lambda_i / x)^\ell \) is between \(-\tau\) and \(\tau\) where \( \tau = \arctan(u/2)/2 \), and we have \( |\arctan(t)| \leq |t| \), so the argument of \( (\lambda_i / x)^\ell \) has smaller absolute value than \( u/4 \), a similar fact holding for \( (\lambda_i / y)^{\ell'} \). As a consequence, there are two integers \( p, p' \) and two numbers \( s \in (-u/4, u/4) \) and \( s' \in (-u'/4, u'/4) \) such that

\[
\ell \omega = p\pi + s \quad \ell' \omega = p'\pi + s'
\]

This implies

\[
p\ell' - p'\ell = \frac{s'\ell - s\ell'}{\pi}.
\]

The LHS is an integer and in the RHS, and the terms \( s'\ell \) and \( s\ell' \) have magnitude smaller than \( u'\ell/4 \) and \( u\ell'/4 \), which were supposed to be smaller than 1 in our hypothesis (8.12), so the whole RHS can only be zero and \( p\ell' = p'\ell \). However, as \( \ell \wedge \ell' = 1 \) we see that \( \ell \) divides \( p \) and \( \ell' \) divides \( p' \), so \( \omega = k\pi + s/\ell \) for some \( k \in \mathbb{Z} \): the complex number \( \lambda \) has argument close to 0 or \( \pi \). We can also see that if \( \theta_i \) is positive, then \( k \) is even and we can indeed take \( \omega = s/\ell \). Otherwise, \( k \) is odd and we can take \( \omega = \pi + s/\ell \).

We may now come back to the relation \( |\lambda_i^\ell - \theta_i^\ell| \leq \varepsilon_0 \), which can also be written as

\[
\lambda_i^\ell = \theta_i^\ell(1 + z)
\]

with \( |z| \leq \varepsilon_0 / |\theta_i|^\ell \). When taking the modulus, we get \( |\lambda_i| = |\theta_i||1 + z|^{1/\ell} \), and from the inequality \( ||1 + z|^{1/\ell} - 1| \leq |z|/\ell \), we finally find out that

\[
||\lambda_i| - |\theta_i|| \leq \frac{\varepsilon_0}{\ell|\theta_i|^{\ell-1}}.
\]

We may finally combining the bounds on arguments and modulus. If \( \theta_i \) is positive, we saw that the argument of \( \lambda_i \) is \( s/\ell \). Writing \( \theta_i = \theta_i e^{is/\ell} + \theta_i - \theta_i e^{is/\ell} \), we find

\[
\square \text{ Springer}
\]}
\[
|\lambda_i - \theta_i| \leq ||\lambda_i| - |\theta_i| + |\theta_i|e^{is/\ell} - 1| \\
\leq (\varepsilon_0/|\theta_i| + |s||\theta_i|)/\ell \\
\leq 2\varepsilon_0/(\ell|\theta_i|^{\ell-1}).
\]

This gives the claimed statement. On the other hand, if \( \theta_i \) is negative, then the argument of \( \lambda_i \) is \( \pi + s/\ell \) and in this case, \( \lambda_i = -|\lambda_i|e^{is/\ell} \) and the same argument holds.

The proof of the statement for the eigenvector is then a consequence of Theorem 8.4 applied to \( A^\ell \).

\section*{9 Proof of Theorem 4.3}

\subsection*{9.1 Notation}

We fix a matrix \( P \) as described above Theorem 4.3.

If \( r_0 \geq 1 \), we define \( \Phi_1 = (\varphi_1, \ldots, \varphi_{r_0}) \) and \( \Sigma = \text{diag}(\mu_1, \ldots, \mu_{r_0}) \). The columns of \( \Phi \) form an orthonormal family, hence \( \Phi^*\Phi = I_{r_0} \).

Recall the parameter \( \ell \) defined in 4.8. The ‘candidate eigenvectors’ are \( u_i = A^\ell \varphi_i / \mu_i^\ell \) and \( v_i = (A^*)^\ell \varphi_i / \mu_i^\ell \), or to put it in matrix form they are the columns of

\[
U = A^\ell \Phi \Sigma^{-\ell} \quad \text{and} \quad V = (A^*)^\ell \Phi \Sigma^{-\ell}.
\]

We set

\[
S = U \Sigma^\ell V^*.
\]

We finally introduce the vector spaces

\[
H = \text{vect}(v_1, \ldots, v_{r_0}) = \text{im}(V) \quad \text{and} \quad H' = \text{vect}(u_1, \ldots, u_{r_0}) = \text{im}(U). (9.1)
\]

Finally, if \( r_0 = 0 \), then \( S \) is simply set to be the zero matrix, and \( H, H' \) are the trivial vector spaces.

\subsection*{9.2 Comments on the Assumptions}

We note that the problem is homogeneous. From now on, we will thus assume without loss of generality that

\[
\mu_1 = 1.
\]

We observe also that the statement is trivial to check if \( \ell = 0 \). In particular, we will assume in the sequel without loss of generality that

\[
\log(n) \geq 8 \log(2d). \quad (9.2)
\]
We will also check easily in 10.3 that $\rho \geq \mu_2^2 = 1$. In particular

$$\vartheta \geq 1 \sqrt{d}$$

It follows that in the statement of Theorem 4.3, we have

$$\tau_0^{-2} \leq d^l \leq n^{1/8}. \quad (9.3)$$

The assumption $C_0 \leq \tau_0^{-2}$ implies that

$$C_0 = cr\tau_0^{4}L^{22} \ln(n)^{16} \leq n^{1/8}. \quad (9.4)$$

### 9.3 Algebraic structure of $A$, $U$, $V$ with respect to $H$

The behaviour of the matrices $U$, $V$ is dictated by a theoretical covariance matrix $\Gamma^{(l)}$, which is a good approximation of the Gram matrices of the columns of $U$ and $V$. It is defined as follows: let $i, j$ be in $[r_0]$ and $t$ be an integer. We will note $\varphi^{i,j}$ for the Hadamard product between $\varphi_i$ and $\varphi_j$:

$$\varphi^{i,j}(x) = \varphi_i(x)\varphi_j(x).$$

Then, we define the matrix $\Gamma^{(l)} \in \mathcal{M}_r(\mathbb{R})$:

$$\Gamma^{(l)}_{i,j} = \sum_{s=0}^{t} \langle 1, Q_s^{*} \varphi^{i,j} \rangle \frac{(\mu_i \mu_j d)^s}{(d_i \mu_j d_j)^s}. \quad (9.5)$$

The diagonal entries $\Gamma^{(l)}_{i,i}$ are exactly the $\gamma_i$ appearing in Theorem 4.3. The next lemma gathers useful properties of $\Gamma^{(l)}$.

**Lemma 9.1** (Properties of $\Gamma^{(l)}$) For any $t$, the matrix $\Gamma^{(l)}$ is a semi-definite positive matrix with eigenvalues greater than 1, and with

$$1 \leq \|\Gamma^{(l)}\| \leq r_0 L^{4} \frac{1 - \tau_0^{2(t+1)}}{1 - \tau_0^2}.$$ 

This lemma will be proved in Sect. 10. The main tool for the subsequent analysis of $U$ and $V$ is the following theorem, which could also be of independent interest.

**Theorem 9.2** (Algebraic structure of $U$ and $V$) There are a universal constant $c > 0$ and an event with probability greater than $1 - cn^{-1/4}$ such that the following holds:

$$\|V^* A^l U - \Sigma^l\| \leq C_1 n^{-1/4} \tau_0^{-2} \rho \ell \quad (9.6)$$

$$\|U^* V - I_{r_0}\| \leq C_1 n^{-1/4} \tau_0^{-2} \rho \ell \quad (9.7)$$
\[ \| \Phi^* U - I_{r_0} \| \leq C_1 n^{-1/4} r_0^{2\ell} \] (9.8)
\[ \| U^* U - \Gamma^{(\ell)} \| \leq C_1 n^{-1/4} r_0^{2\ell} \] (9.9)
\[ \| V^* V - \Gamma^{(\ell)} \| \leq C_1 n^{-1/4} r_0^{2\ell} \] (9.10)
\[ \| A^\ell P_{H^\perp} \| \leq C_2 \delta^\ell \] (9.11)
\[ \| P_{H^\perp} A^\ell \| \leq C_2 \delta^\ell , \] (9.12)

where \( C_1 = c r_0 b^2 \ln(n)^{5/2} \) and \( C_2 = c L^{10} r \ln(n)^{12} \).

**Remark 9.3** The same statements holds, with the same constants, if we replace \( \ell \) by \( \ell' = \ell + 1 \). Note that \( \ell' \) and \( \ell \) are then mutually prime.

The proof of this theorem occupies the next sections of this paper. We now use this theorem to prove all the results mentioned before.

### 9.4 Proof of Theorem 4.3

We consider the intersection of the two events of Theorem 9.2 for \( \ell \) and \( \ell' = \ell + 1 \) which has probability at least \( 1 - cn^{-1/4} \). Let us call \( \mathcal{E}_n \) this event. Our goal is to apply Theorem 8.6 to \( A^\ell \) and \( S \) and to \( A^{\ell'} \) and \( S' \) on \( \mathcal{E}_n \). Our main task will be to check the three conditions of Theorem 8.6; we will focus on checking them for \( U \) and \( V \), the statements for \( U' \), \( V' \) working obviously in the same way.

We will use that 9.4 implies that \( C_1 n^{-1/4} \) goes to 0 as \( n \) goes to infinity, uniformly in \( r, b \) satisfying (9.4). In particular, for all \( n \geq n_0 \) large enough

\[ C_1 n^{-1/4} \leq 1/2. \] (9.13)

In the sequel, we always assume that \( \mathcal{E}_n \) holds and that \( n \geq n_0 \).

**Condition 1 of Theorem 8.6** This is settled by Remark 9.3.

**Condition 3 of Theorem 8.6** From (9.9) and Lemma 9.1, we have,

\[ \| U \|^2 = \| U^* U \| \leq (\ell + 1) r_0 L^4 + C_1 n^{-1/4} \leq r_0 L^4 \ln(n) \] (9.14)

the last line coming from 9.13. The same inequality holds for \( \| V \|^2 \). We find that

\[ \| U \| \vee \| V \| \leq N := \sqrt{r_0 L^4 \ln(n)}. \]

For the conditioning properties of \( U, V \), we deduce from (9.13) and Lemma 9.1 that

\[ \| U^* U \| \wedge \| V^* V \| \geq \lambda_{\min} (\Gamma^{(\ell)}) - C_1 n^{0.1} \geq h := 1/2. \] (9.15)

Similarly, from 9.7, we have

\[ \| U^* V - I_{r_0} \| \leq C_1 n^{-1/4} r_0^{2\ell} := \delta. \] (9.16)
This gives condition 3 for $U$, $V$. We note that (9.15) implies that:

$$
\| \left( U^* U \right)^{-1} \| \vee \| \left( V^* V \right)^{-1} \| \leq 2.
$$

(9.17)

**Condition 2 of Theorem 8.6** This condition requires more work. We have

$$
P_H = V \left( V^* V \right)^{-1} V^* \text{ and } P_{H'} = U \left( U^* U \right)^{-1} U^*.
$$

We also note that $P_{H \perp} = I_n - P_H$ the projection matrix on $H \perp$. Since $SP_{H \perp} = P_{H \perp} S = 0$, we have $S = P_{H \perp} SP_H$ and

$$
\| A^\ell - S \| \leq \| P_{H'} A^\ell P_{H} - S \| + \| A^\ell P_{H \perp} \| + \| P_{H'} A^\ell P_{H} \|
\leq \| U \| \| \left( U^* U \right)^{-1} U^* A^\ell V \left( V^* V \right)^{-1} - D^\ell \| \| V \| + \| A^\ell P_{H \perp} \| + \| P_{H'} A^\ell P_{H} \|,
$$

(9.18)

where at the second line, we have used that $U D^\ell V^* = S$. To bound the above expression, we will use that the fact that $V^* U - I_{r_0}$ is small implies that $U$ is close to $\tilde{U} := V \left( V^* V \right)^{-1}$ and that $V$ is close to $\tilde{V} := U \left( U^* U \right)^{-1}$. More precisely, we write

$$
P_H U = \tilde{U} V^* U
= \tilde{U} + E_1
$$

where $E_1 := \tilde{U} \left( V^* U - I_{r_0} \right)$. From (9.14–9.16–9.17), we find that

$$
\| E_1 \| \leq 2 \delta N.
$$

We may thus decompose $U$ as follows:

$$
U = \tilde{U} + E_1 + P_{H \perp} U.
$$

Similarly, we find

$$
V = \tilde{V} + E_2 + P_{H' \perp} V
$$

with $E_2 = \tilde{V} \left( U^* V - I_{r_0} \right)$ and $\| E_2 \| \leq 2 \delta N$. We get the following:

$$
\| \tilde{V}^* A^\ell \tilde{U} - \Sigma^\ell \| \leq \| \tilde{V}^* A^\ell U - \Sigma^\ell \| + \| \tilde{V}^* A^\ell P_{H \perp} \| + \| \tilde{V} \| \| A^\ell \| \| E_1 \| + \| A^\ell P_{H \perp} \| \| U \|
\leq \| V^* A^\ell U - \Sigma^\ell \| + \| \tilde{V}^* A^\ell \| \| E_1 \| + \| A^\ell P_{H \perp} \| \| U \|
+ \| U \| \| A^\ell \| \| E_1 \| + \| P_{H' \perp} A^\ell \| \| V \|.
$$

We use this last inequality in (9.18) and use (9.6). We obtain the bound:

$$
\| A^\ell - S \| \leq 2 N^2 \delta^\ell + 4 N^3 C_2 \delta^\ell + 4 \delta N^4 \| A^\ell \|.
$$
From our choice of parameters 9.4, we have

\[ 4\delta N^4 \leq 1/2. \]  

(9.19)

It follows that

\[ \|A^\ell - S\| \leq 5N^3C_2\theta^\ell + 4\delta N^4\|A^\ell\|. \]  

(9.20)

We may use 9.20 and \( \|S\| \leq \|U\|\Sigma\|V\| \leq N^2 \) to upper bound \( \|A^\ell\| \). Indeed we write \( \|A^\ell\| \leq \|S\| + \|A^\ell - S\| \leq N^2 + \|A^\ell - S\| \), we find from 9.20 and 9.19:

\[ \|A^\ell\| \leq 2N^2 + 10N^3C_2\delta^\ell. \]  

(9.21)

Putting this last expression back in 9.20) and using \( 4\delta N^4 \leq 1/2 \), we obtain the bound:

\[ \|A^\ell - S\| \leq \eta := 10N^3C_2\theta^\ell + 8\delta N^6. \]  

(9.22)

The conclusion of Theorem 4.3 is then a consequence of Theorem 8.6. Let us now estimate roughly the quantity \( \epsilon \) in Theorem 8.6. We first note that \( 2\delta N^6 \) is larger than \( \delta N^{4}\sqrt{r_0}/h \), since \( h = 1/2 \) and \( N\sqrt{r_0} \leq N^2 \). It follows that \( \epsilon \) in Theorem 8.6 is bounded by \( 24\eta N^3 \). We claim also that

\[ \delta N^6 \leq N^3C_2\theta^\ell. \]  

(9.23)

Indeed, \( n^{-1/4} = n^{-\kappa} \leq (2d)^{-2\ell} \) and, since \( \rho \geq \mu_1 \geq 1 \) (see forthcoming bound (10.3)), we have \( \theta \geq 1/\sqrt{d} \). It follows that \( n^{-1/4} \leq \theta^\ell \). To prove (9.23), we then need to check that \( \tau_0^{2\ell}C_1N^3 \leq C_2 \), the latter is immediate using 4.9.

It follows that, for some universal constant \( c, c' > 0 \),

\[ \epsilon \leq cN^6C_2\theta^\ell \leq c'rr_0^3L^{12}L^{10} \ln(n)^{15}\theta^\ell. \]

We set \( C_0 = 4\epsilon r_0 \ln(n)/\theta^\ell \). After crude rearrangement and simplifications, the statement of Theorem 4.3 for eigenvalues follows easily from Theorem 8.6 applied to \( r = r_0 \).

For eigenvectors, Theorem 8.6 allows us to describe the behaviour of the eigenvectors of \( A \). We may assume without loss of generality that

\[ \frac{C_0^{2\ell}}{(1 - \tau_{i,\ell})^2} \leq 1, \]

since otherwise the statement is trivial (in particular \( \tau_{i,\ell} < 1 \)). In particular, for any \( i \in [r_0] \), we have

\[ (1 - \tau_{i,\ell})\theta^\ell \geq (1 - \tau_{i,\ell})\tau_0^{-\ell}\theta^\ell \geq C_0\theta^\ell \geq 4\epsilon. \]

We are thus in position to apply the eigenvector part in Theorem 8.6.
Note that \(|\theta_i|^\ell - 2\varepsilon \geq |\theta_i|^\ell /2\) and thus \(N^2 \eta /|\theta_i|^\ell - 2\varepsilon \leq 2N^2 \eta \leq \varepsilon / (6N)\). We choose the orientation of \(\psi_i\) to ensure that \(\langle \psi_i, u_i \rangle\) is non-negative. We deduce from 8.14 that,

\[
\left| \psi_i - \frac{u_i}{|u_i|} \right|^2 = 2 - 2\langle \psi_i, \frac{u_i}{|u_i|} \rangle \leq \frac{40\varepsilon^2}{(1 - \tau_{i,\ell})^2|\theta_{r_0}|^2} \leq \frac{40C_0^2 \tau_{0,\ell}^2}{(1 - \tau_{i,\ell})^2}.
\]

We then write

\[
\left| \langle \psi_i, \varphi_j \rangle - \frac{u_i}{|u_i|}, \varphi_j \right| \leq \left| \psi_i - \frac{u_i}{|u_i|} \right|.
\]

Thanks to (9.8), we have \(\langle u_i/|u_i|, \varphi_j \rangle\) is close to \(\delta_{i,j}/|u_i|\) with error bounded by \(\delta/|u_i|\). Finally, 9.9 implies that

\[
\left| \frac{|u_i|^2 - \gamma_i}{|u_i| + \sqrt{\gamma_i}} \right| \leq \frac{\delta}{\sqrt{\gamma_i}}.
\]

When gathering all those bounds, we get that \(\langle \psi_i, \varphi_j \rangle - \delta_{i,j}/\gamma_i\) is bounded by \(cC_0 \tau_{0,\ell}^2/(1 - \tau_{i,\ell})\) where \(c\) is a universal constant. Defining a new constant \(C_0\) equal to \(cC_0\), we thus obtain the bound displayed in 4.12. The same proof with 9.7 in place of 9.8 gives the bound 4.16. It concludes the proof of Theorem 4.3.

### 9.5 Proof of Corollary 4.6

On the event of Theorem 4.3, we have

\[
\left| \langle \psi_i + \psi_i', \varphi_j \rangle - \frac{2}{\sqrt{\gamma_i}} \frac{\delta_{i,j}}{|\psi_i + \psi_i'|} \right| \leq \frac{2C_0 \tau_{0,\ell}^2}{1 - \tau_{i,\ell}} =: 2\varepsilon.
\]

Moreover, since \(|\psi_i + \psi_i'|^2 = 2 + 2\langle \psi_i, \psi_i' \rangle\), we have

\[
||\psi_i + \psi_i'|^2 - 2 - 2/\gamma_i| \leq 2\varepsilon.
\]

In particular, since \(\gamma_i \geq 1\),

\[
||\psi_i + \psi_i'| - \sqrt{2 + 2/\gamma_i}| \leq 2\varepsilon/(|\psi_i + \psi_i'| + \sqrt{2 + 2/\gamma_i}) \leq \sqrt{2\varepsilon}.
\]
If $\varepsilon \leq 1/2$, then $|\psi_i + \psi'_i| \geq 1$ and we find

$$\left| \frac{1}{|\psi_i + \psi'_i|} - \frac{1}{\sqrt{2 + 2/\gamma_i}} \right| \leq \frac{\sqrt{2}\varepsilon}{\sqrt{2 + 2/\gamma_i}} \leq \varepsilon.$$ 

We thus have checked that if $\varepsilon \leq 1/2$ then

$$\left| \frac{\langle \psi_i + \psi'_i | \psi_i + \psi'_i \rangle - \sqrt{2}}{\sqrt{\gamma_i + 1}} \delta_{i,j} \right| \leq 4\varepsilon.$$ 

Otherwise, $\varepsilon > 1/2$ and this last inequality also holds since it is trivial in this case.

### 10 Consequences of Algebraic Incoherence

Our goal in this section is to gather several useful estimates on $Q$ and $\Gamma^{(1)}$ linked with the incoherence properties 4.2 and 4.3.

#### 10.1 Incoherence of $P$ and $Q$

The parameter $L$ in (4.2) is not independent of the parameter $b$ in 4.3.

We introduce the scale invariant analog of the parameter $L$ for the matrix $Q$: we set

$$K = n \max_{x,y} Q_{xy}/\rho.$$  

(10.1)

We have

$$L = n \max_{x,y} |P_{xy}| = \sqrt{K\rho}.$$  

Notice also that the scalar $K$ is scale invariant. We note also that the following bound holds:

$$1 \leq K \leq \frac{L^2}{\mu_1^2}. $$  

(10.2)

Indeed, for the lower bound, we use that for any matrix $T$, $\|T\| \leq n \max_{x,y} |T_{xy}|$. For the upper bound, we use that $K = L^2/\rho^2$ and

$$\rho \geq \frac{\langle 1, Q1 \rangle}{n} = \sum_{x,y} P_{x,y}^2 = \|P\|_F^2 \geq \mu_1^2.$$  

(10.3)
10.2 Bounds on the Entries of $Q$

We start by bounding the entries of powers of $Q$. For any $x$, we find

$$\sum_y Q_{xy} \leq K \rho.$$ 

It follows that for any $x, y$,

$$(Q^2)_{xy} = \sum_z Q_{xz} Q_{zy} \leq \frac{K^2 \rho^2}{n}.$$ 

Let $(\psi_k)$ be an ON basis of eigenvectors of $Q$ with eigenvalues $(\nu_k)$. Let $t \geq 2$, we write for any $x, y$,

$$(Q^t)_{xy} = \sum_k \nu_k^t \psi_k(x) \psi_k(y) \leq \rho^{t-2} \sum_k \nu_k^2 |\psi_k(x)||\psi_k(y)| \leq \rho^{t-2} \sqrt{(Q^2)_{xx}} \sqrt{(Q^2)_{yy}},$$

where the last step follows from Cauchy-Schwarz inequality. In particular, for any $t \geq 2$ and $x, y$

$$(Q^t)_{xy} \leq \frac{K^2 \rho^t}{n} \quad (10.4)$$

It follows from 10.1 that Eq. (10.4) also holds for $t = 1$. The following immediate consequence will be crucial; the idea it conveys is that, for any vector $v, w$, $\langle v, Q^t w \rangle$ is essentially bounded by $|v|_1 |w|_1 \rho^t / n$, a result in the flavour of Perron-Frobenius theory.

Lemma 10.1 For any integer $t \geq 1$ and any vectors $v, w \in \mathbb{R}^n$,

$$\langle v, Q^t w \rangle \leq \frac{|v|_1 |w|_1 K^2 \rho^t}{n}.$$ 

Proof We simply write

$$\langle v, Q^t w \rangle = \sum_{x,y} (Q^t)_{x,y} v(y) w(y) \leq \sum_{x,y} \frac{K^2 \rho^t}{n} |v(x)||w(y)|$$

where we have used 10.4. The conclusion follows. \qed
10.3 Proof of Lemma 9.1: The Incoherence Hypothesis for the Covariance Matrix

Let us end this section by the proof of Lemma 9.1. We start by recalling the definition of the theoretical covariance $\Gamma^{(t)}_{i,j} \in \mathcal{M}_{r_0, r_0}(\mathbb{R})$:

$$
\Gamma^{(t)}_{i,j} = \sum_{s=0}^{t} \frac{\langle 1, Q^s \varphi^{i,j} \rangle}{(\mu_i \mu_j d)^s}.
$$

By Lemma 10.1, we find

$$
\langle 1, Q^s \varphi^{i,j} \rangle \leq |\varphi^{i,j}|_1 \leq K^2 \rho^t \leq K^2 \rho^t (10.5)
$$

where we have use Cauchy-Schwarz inequality:

$$
|\varphi^{i,j}|_1 = \sum_x |\varphi_i(x)||\varphi_j(x)| \leq |\varphi_i|_2|\varphi_j|_2 = 1.
$$

Going back to the sum defining $\Gamma^{(t)}_{i,j}$, we get

$$
|\Gamma^{(t)}_{i,j}| \leq K^2 \frac{1 - (\rho/(\mu_i \mu_j d))^{t+1}}{1 - \rho/(\mu_i \mu_j d)} \leq \frac{K^2 (1 - \tau_0^{2(t+1)})}{1 - \tau_0^2}
$$

and as a consequence,

$$
\|\Gamma^{(t)}\| \leq r_0 \max_{i,j} |\Gamma^{(t)}_{i,j}| = \frac{r_0 K^2 (1 - \tau_0^{2(t+1)})}{1 - \tau_0^2}
$$

On the other hand, if we note

$$
C^{(s)}_{i,j} = \frac{\langle 1, Q^s \varphi^{i,j} \rangle}{(\mu_i \mu_j d)^s}
$$

then it is not difficult to see that $C^{(s)}$ is indeed a semi-definite positive (SDP) matrix; more precisely, if we introduce $\pi_s(x) = \sqrt{Q^s \mathbf{1}(x)} \geq 0$ and $\Pi_s = \text{diag}(\pi_s)$, then

$$
C^{(s)} = d^{-s} \cdot D^{-s} \Phi^* \Pi_s^2 \Phi D^{-s} \quad (10.6)
$$

which is clearly SDP. The matrix $\Gamma^{(t)}$ is thus a sum of $C^{(0)} = I_{r_0}$ and $t - 1$ SDP matrices, hence it is itself an SDP matrix and its eigenvalues are greater than the eigenvalues of $I_{r_0}$, hence the first statement.

For further needs, we notice that if $i \in [r] \setminus [r_0]$, then $|\mu_i| \leq \vartheta$ and we get from 10.5 that

$$
\Gamma^{(t)}_{i,i} \leq \sum_{s=0}^{t} \frac{K^2 \vartheta^{2s}}{|\mu_i|^{2s}} \leq K^2 (t + 1) \frac{\vartheta^{2t}}{|\mu_i|^{2t}}. \quad (10.7)
$$
11 Coupling Graphs and Trees

The basic ingredients for the proofs of Theorem 4.3 and related statements are directed Galton-Watson trees and martingales defined on them. We start to introduce the notations and vocabulary for this.

11.1 Marked Galton-Watson Trees and Erdos-Rényi Digraphs

11.1.1 Graph-Theoretic Definitions

A marked graph with mark space \( \mathbb{N} \) is a digraph \((V, E)\), with possible loops, endowed with a mark function \( t : V \to \mathbb{N} \). We are going to note \( \mathcal{G}_s \) the set of all rooted directed graphs on a common countable set \( V \) and with mark space \( \mathbb{N} \). Formally, the elements of \( \mathcal{G}_s \) are triples \((G, o, t)\), with \( o \) the root, but in general we will drop the mark function \( t \) and simply write \((G, o)\).

Let \((G, o, t) \in \mathcal{G}_s\) and \( g = (V, E) \). If \( W \subset V \) is a subset of \( V \) containing the root, then the induced subgraph \((G, o, t)_W\) is defined as follows: the underlying graph is \( G_W := (V, E_W) \) where \((i, j) \in E_W \) if and only if \((i, j) \in E \) and both \( i \) and \( j \) are in \( W \), and the mark function \( t_W \) is given by \( t_W(v) = t(v) \) for all \( v \in W \).

The elements in \( \mathcal{G}_s \) are digraphs, and therefore we need to make a distinction between directed paths and undirected paths. Let \( g = (V, E) \) be a digraph,

- If \((x, y) \in E\) we note \( x \rightarrow y \),
- if \( x \rightarrow y \) or \( y \rightarrow x \) or both, we note \( x \sim y \).

Every directed graph \( G \) can be transformed into an undirected graph \( \hat{G} = (V, \hat{E}) \) by simply forgetting the direction of the edges: \((x, y) \in \hat{E} \) iff \( x \sim y \) in \( g \).

If \( u, v \in V \), a directed path or dipath from \( x \) to \( y \) is a sequence of vertices \( x_0 = x, x_1, \ldots, x_k = y \) such that for every \( s \) we have \( u_s \to u_{s+1} \). A path is the same except that we only ask \( x_s \sim x_{s+1} \).

- The length of the shortest directed path between \( x \) and \( y \) is denoted by \( d^+(x, y) \).
- The length of the shortest directed path between \( y \) and \( x \) is also denoted by \( d^-(x, y) = d^+(y, x) \).
- The length of the shortest path is denoted by \( d(x, y) \).
- When \( G \) is a digraph graph, \( \mathcal{P}_G(x, t) \) is the set of paths in \( G \) starting from \( x \) and having \( t \) steps.

The set of all \( y \) such that \( d^+(x, y) \leq t \) is the forward ball \( B^+_G(x, t) \) and the set of all \( y \) such that \( d(x, y) \leq t \) is the ball \( B_G(x, t) \). When no confusion can arise, we write \( B^+ \) or \( B \) instead of \( B_G, B^+_G \). If \( t \) is an integer and \((G, x) \in \mathcal{G}_s \), then \((G, x)_t \) is the subgraph of \((G, x)\) induced by \( B_G(x, t) \), as defined above, and similarly \((G, x)_t^+ \) is the subgraph of \((G, x)\) induced by \( B^+(x, t) \).

A cycle in the graph \( G \) is a sequence of distinct vertices \((x_1, \ldots, x_k)\) such that \( x_s \sim x_{s+1} \) for every \( s < k \) and \( x_k \sim x_1 \). The number \( k \) is the length of the cycle.

A tangle-free subgraph of \( G \) is a subgraph of \( G \) that contains at most one cycle. The graph \( G \) is \( t \)-tangle free if for every vertex \( x \), the ball \( B_G(x, t) \) is tangle-free.
We will the following useful property of tangle freeness. If \( G \) is \( h \)-tangle free then there is either zero or one cycle in \((G, x)\), for \( t \leq h \). Hence for any \( y \) at distance \( t \) from \( x \), there is at most two paths of length \( t \) from \( x \) to \( y \). To put it another way, for any \( x \in [n] \), we have for \( t \leq h \),
\[
|\mathcal{P}_G(x, t)| \leq 2|((G, x)_t)|. \tag{11.1}
\]

### 11.1.2 Definition of the Graph \( G \) and the Marked Galton Watson Tree

We define that \( G \) as the directed Erdős-Rényi graph (with loops) whose adjacency matrix is given by \( M \): the vertex set is \([n]\) and each directed edge is present independently with probability \( d/n \). Let \( x \) be an arbitrary element of \([n]\). We root the graph \( G \) at \( x \), and we mark every vertex with itself: the mark of vertex \( x \in [n] \) is simply the integer \( \iota(x) = x \). The resulted marked graph \((G, x)\) is an element of \( \mathcal{G}_s \).

We now define the directed Galton-Watson tree \( T \) in the following way. Starting from its root \( o \), every vertex has a \( \text{Poi}(2d) \) number of children. Every edge \((u, v)\) is independently given a unique direction \( u \rightarrow v \) or \( v \rightarrow u \) with probability 1/2. This yields a random directed tree. Equivalently, each vertex has a \( \text{Poi}(d) \) number of ‘out-children’ and a \( \text{Poi}(d) \) number of ‘in-children’.

Finally, every non-root vertex \( o' \) is independently given a random mark \( \iota(o') \) which is uniform on \([n]\). The root is given a special mark \( \iota(o) = \mathcal{O} \). The resulting element of \( \mathcal{G}_s \) should be noted \((T_n, \mathcal{O})\) because it depends on \( n \) through the marks, but we will simply note \((T, \mathcal{O})\). We shall say that the tree \((T, \mathcal{O})\) is **grown from the seed** \( \mathcal{O} \).

### 11.2 Growth Properties: Trees

Let us first state several properties on the growth of the tree \( T \) first, then on the graph \( G \). They are directly drawn from \([14]\), see Section 8 for the tree, and Sections 9.1-9.2 for the graph. Clearly, the underlying undirected tree obtained from \( T \) by deleting the marks and orientations is simply a \( \text{Poi}(2d) \) Galton-Watson tree, which allows us to use known results on the growth properties of GW trees. We recall \( D = 2d \lor 1.01 \) was defined in Theorem 4.3.

**Lemma 11.1** ([14, Lemma 23]) Let us note \( S_t \) the number of vertices at distance \( t \) from the root \( o \) of \( T \). There are two universal constants \( c_0, c_1 > 0 \) such that for all \( \lambda > 0 \),
\[
P(S_t \leq \lambda D^t \text{ for all } t) \geq 1 - c_0 e^{-c_1 \lambda}. \tag{11.2}
\]
Moreover, there is a universal constant \( c \) such that for every \( p \geq 1 \),
\[
E \left[ \max_{t \geq 1} \left( \frac{S_t}{D^t} \right)^p \right] \leq (cp)^p. \tag{11.3}
\]
Proof. The second claim is an immediate consequence of the first: we have

\[
E \left[ \max_{t \geq 1} \left( \frac{S_t}{(2d)^t} \right)^p \right] = p \int_0^\infty \lambda^{p-1} \mathbf{P} \left( \max_{t \geq 1} \left( \frac{S_t}{D^t} \right) \geq \lambda \right) d\lambda \leq p \int_0^\infty \lambda^{p-1} c_0 e^{-c_1 \lambda} d\lambda.
\]

The first statement of the lemma is [14, Lemma 23], it is however not explicitly written in the proof of [14, Lemma 23] that the constants \( c_0, c_1 \) are universal. An inspection of the proof shows this is indeed the case (we use here that \( D \) is bounded away from 1).

As a consequence we may easily upper bound the size of \((T, \mathcal{O})_t\). Indeed, we use the inequality

\[
|(T, \mathcal{O})_t| = 1 + S_1 + \cdots + S_t \\
\leq \max_k (D^{-k} S_k) \sum_{k=0}^t D^k \\
\leq 2 \max_k (D^{-k} S_k) D^t,
\]

where we have use that \((D^{t+1} - 1)/(D - 1) \leq 2D^t\) for all \( D \geq 2 \). Taking expectation, we deduce from Lemma 11.1 that for any integers \( p \geq 1 \) and \( t \geq 1 \),

\[
\left( E[|(T, \mathcal{O})_t|^p] \right)^{1/p} \leq 2c_p D^t. \tag{11.4}
\]

11.3 Growth Properties: Graphs

We now establish the same properties as before, but for the directed graph \( G \) whose adjacency matrix is \( M \). We start by proving that up to a depth of order \( \log_2 d n \), the graph \( G \) has few cycles. We denote by \( \hat{G} \) the graph \( G \) in which the directions have been erased. Pick any vertices \( x, y \). Then, from the union bound

\[
P([x, y] \in E(\hat{G})) = P((x, y) \in E(G) \text{ or } (y, x) \in E(G)) \leq \frac{D}{n},
\]

where as above we have set \( D = 2d \). This shows that the edge distribution of \( \hat{G} \) is stochastically dominated by the edge distribution of an undirected \((n, D/n)\) Erdős-Rényi graph (note that, unlike the usual definition, here we allow loops in the random graph). Events which are monotone for the deletion of edges (such has having a few number of cycles) are thus of smaller probability in \( G \) than in \( \text{ER}(n, D/n) \). As a consequence some results in [14] directly transfer to our setting.

Lemma 11.2 (Growth rate, [14, Lemma 29]) Let us denote by \( S_t(x) \) the number of vertices in \( G \) that are exactly at (unoriented) distance \( t \) from vertex \( x \). There are two universal constants \( c_0, c_1 > 0 \) such that for every positive \( \lambda \) and every vertex \( x \in [n] \), we have
\begin{equation}
\mathbf{P}(S_t(x) \leq \lambda D^t \text{ for all } t) \geq 1 - c_1 e^{-c_0 \lambda}.
\end{equation}

Moreover, there is a universal constant \( c \) such that for every \( p \geq 1 \),
\begin{equation}
\mathbb{E} \left[ \max_{t \geq 1} \left( \frac{S_t(x)}{D^t} \right)^p \right] \leq (cp)^p \quad \text{and} \quad \mathbb{E} \left[ \max_{x \in [n]} \left( \frac{S_t(x)}{D^t} \right)^p \right] \leq (c \ln n)^p + (cp)^p.
\end{equation}

**Proof** Only the LHS inequality 11.6 is not explicitly stated in [14]. It is proved as in the proof of Lemma 11.1. \( \square \)

Let us apply this result to \(|(G, x)_t| = 1 + S_1(x) + \cdots + S_t(x)\) with \( \lambda := c^{-1}_0 \ln(c_1 n^2) \).

With probability greater than \( 1 - 1/n \), for any \( t \) and for any \( x \in [n] \),
\begin{align*}
S_t(x) &\leq c \ln(n) D^t
\end{align*}

where \( c \) is a universal constant. On this event, one also has
\begin{align*}
|(G, x)_t| &= 1 + S_1(x) + \cdots + S_t(x) \\
&\leq c \ln(n)(1 + D + \cdots + D^t) \\
&\leq 2c \ln(n) D^t.
\end{align*}

Similarly, we use that
\begin{align*}
|(G, x)_t| &\leq 2D^t \max_{k \geq 0} \frac{S_k}{D^k}
\end{align*}

(as explained above 11.4) and we deduce from 11.6 that for any \( p \geq 1 \),
\begin{align*}
\mathbb{E} \left[ |(G, x)_t|^p \right]^{1/p} &\leq 2cpD^t. \\
\mathbb{E} \left[ \max_{x \in [n]} |(G, x)_t|^p \right]^{1/p} &\leq 2c(\ln(n) + p)D^t.
\end{align*}

**11.4 Distance Between Neighborhoods in the Graph and the GW Tree**

In this section, we fix some \( \kappa \geq 0 \) and we consider an integer \( h \) such that
\begin{equation}
0 \leq h \leq \kappa \log_D(n)
\end{equation}

where we have set \( D = 2d \lor 1.01 \) as above.

Our goal is to study fine geometric properties of \((G, x)_h\) for \( \kappa \) small enough. The above comparison trick between \( G \) and undirected Erdős-Rényi with parameters \((n, D/n)\) implies that the following holds:
Lemma 11.3 (Tangle-free, [14, Lemma 30]) Let \(0 \leq \kappa \leq 0.49\) and \(h\) an integer as in (11.10). For some universal constant \(c\), the graph \(G\) is \(h\)-tangle free, with probability at least \(1 - cn^{2\kappa - 1}\). Moreover, for any vertex \(x \in [n]\), the graph \((G, x)_h\) has no cycle with probability greater than \(1 - cn^{\kappa - 1}\).

**Proof** In [14, Lemma 30] there are no loops. The probability of having a loop at vertex \(x\) is \(\frac{d}{n}\). It is however immediate to check that the same proof extends also in our case. \(\square\)

We quantify the distance between neighborhoods of \(G\) and \(T\) up to the depth \(h\).

Let us recall some definitions. If \(P_1, P_2\) are two probability measures on the space \((\Omega, F)\), their total variation distance is defined as

\[
d_{TV}(P_1, P_2) = \min_{(X_1, X_2) \in \pi(P_1, P_2)} P(X_1 \neq X_2)
\]

where \(\pi(P_1, P_2)\) denotes the set of couplings between \(P_1\) and \(P_2\): pairs of random variables \((X_1, X_2)\) such that \(X_1\) is distributed as \(P_1\) and \(X_2\) is distributed as \(P_2\). It is a well-known fact (see [42]) that the total variation distance is also given by

\[
d_{TV}(P_1, P_2) = \max_{A \in F} P_1(A) - P_2(A).
\]

We note \(\mathcal{L}(X)\) the probability distribution of a random variable \(X\).

**Proposition 11.4** (GW-tree approximation) Let \(0 \leq \kappa \leq 0.49\) and \(h\) an integer as in (11.10). There is a universal constant \(c > 0\) such that for every vertex \(x\),

\[
d_{TV}(\mathcal{L}((G, x)_h), \mathcal{L}((T, x)_h)) \leq c(\ln n)^2 n^{2\kappa - 1}.
\]  

The proof of this fact is classical; one can adapt the arguments in [14] to our setting. The difference is that our graphs are directed and now have \([n]\) possible labels, but this only brings shallow difficulties. We sketch the main ideas.

Let us recall the following very classical total variation distance:

\[
d_{TV}(\text{Bin}(n, \lambda/n), \text{Poi}(\lambda)) \leq \frac{\lambda}{n}.
\]  

**Coupling between labelled graphs**

As a consequence of Lemma 11.3 and (11.7), with a probability greater than \(1 - cn^{\kappa - 1}\), the graph \((G, x)_h\) is a directed tree and contains no more than \(k := c \ln(n)n^\kappa\) vertices. Let us note \(E_h\) this event and perform a breadth-first exploration starting from \(x\). This exploration finishes at a time \(\tau \leq k\). At each step, we reveal a set of \(\text{Poi}(d)\) out-vertices and \(\text{Poi}(d)\) in-vertices. From 11.12, we make a total-variation error smaller than \(d/(n-k) + d/(n-k) \leq 3d/n\) for \(n\) large enough. By repeatedly conditioning, the total variation error made on \(E_h\) is not greater than \(3d\tau/n \leq 3dk/n\). This gives
a coupling between the unlabelled versions of \((G, x)_h\) and \((T, x)_h\) which fails with probability at most \((c + 3cd) \ln(n)n^{k-1} \leq 4cd \ln(n)n^{k-1} \).

We now bring the labels in. With probability greater than \(4cd \ln(n)n^{k-1}\), the coupling between the unlabelled versions of \((G, x)_h\) and \((T, x)_h\) succeeds and have size smaller than \(k = c \ln(n)n^k\). We then put the labels in the Erdős-Rényi graph by drawing a uniform ordered \(k\)-set from \([n]\), while we put the labels on the Galton-Watson tree by simply drawing \(k\) i.i.d. uniform samples from \([n]\). We claim that the total variation distance between these two random multi-sets is smaller than \(k^2/n\), see below for the proof. Hence the labels agree with an extra total variation cost of \((c \ln(n)n^k)^2/n\).

In the end, the coupling created this way fails with probability at most \(4cd \ln(n)n^k + c^2 (\ln(n))^2 n^{2k-1}\) which is exactly what is needed, up to adjusting the constants (note that we may assume that \(D \leq n^k\) otherwise \(h = 0\) and the statement is trivial). This concludes the sketch of proof of Proposition 11.4 up to the claimed bound \(k^2/n\) for the distance between the random multisets which we now explain.

**Sampling with and without replacement**

Let \(m\) be an integer. We define two random multisets in the following way. Put \(m\) identical balls with labels from 1 to \(m\) in a big urn. Draw the first ball and set \(p_1\) and \(q_1\) to be its label. Put the ball back in the urn. Then, suppose that one has constructed \((p_1, \ldots, p_t)\) and \((q_1, \ldots, q_t)\). Do the following:

- Draw a ball from the urn and set \(p_{t+1}\) to be the label of this ball.
- If this label is not already one of the \(q_s\), set it onto \(q_{t+1}\). Else, put the ball back in the urn and draw as many balls as needed to get a label which is not already one of the \(q_s\). Define \(q_{t+1}\) to be this label.

It is clear that for every \(k \leq m\), \(Q_k := (q_1, \ldots, q_k)\) is a uniform ordered \(k\)-set from \([m]\), while \(P_k := (p_1, \ldots, p_k)\) is distributed as \(k\) i.i.d. uniform elements in \([m]\). The random variable \((P_k, Q_k)\) is thus a coupling between those two distributions. This coupling is successful if and only if \(P_k\) has exactly \(k\) distinct elements, which happens with probability

\[
\frac{(m-1) \ldots (m-k-1)}{m^k} \geq \left( 1 - \frac{k}{m} \right)^k \geq 1 - \frac{k^2}{m}.
\]

The coupling thus fails with probability smaller than \(k^2/m\), an upper bound for the total-variation distance between \(P_k\) and \(Q_k\).

**Maximal coupling of trees and graphs**

Proposition 11.4 tells us that for every fixed \(x\), there exists a random rooted marked tree \((T_x, x)\) defined on the same probabilistic space as \((G, x)\) and such that

\[
P((T_x, x)_h \neq (G, x)_h) \leq c (\ln n)^2 n^{2k-1}.
\]

provided that \(h \leq \kappa \log_D n\) with \(\kappa \leq 0.49\) and \(D = 2d\). In the sequel, we will use this family of coupled trees \((T_x, x)\) for \(x \in [n]\).
12 Graph Functionals

12.1 Functionals on Trees: Computations

We now introduce a family of functionals on $G^*$ that will be used several times in the sequel. Remember that when $(g, o)$ is a rooted marked graph, we note $P_g(o, t)$ the number of paths in $g$ starting from the root $o$ and having $t$ steps, that is, $(t + 1)$-uples $x = (x_0, x_1, \ldots, x_t)$ with $x_0 = o$ and $x_s \rightarrow x_{s+1}$.

In this section, $\psi, \phi$ represent two vectors in $\mathbb{R}^n$ and $t$ is an integer. We define

$$f_{\phi, \psi, t}(g, o) = \left(\frac{n}{d}\right)^t \phi(t(o)) \sum_{P_g(o, t)} P_{t_1(x_1)} \cdots P_{t(x_t)} \times \psi(t(x_t)).$$ \hspace{1cm} (12.1)

We clearly have

$$f_{\phi, \psi, t}(G, x) = \phi(x)(A^t\psi)(x).$$ \hspace{1cm} (12.2)

We will also need another functional:

$$F_{\mu, \psi, t}(g, o) = f_{1, \psi, t}(g, o) - \mu^{-1} f_{1, \psi, t+1}(g, o).$$ \hspace{1cm} (12.3)

We have

$$F_{\mu, \psi, t}(G, x) = A^t\psi(x) - \frac{1}{\mu} A^{t+1}\psi(x).$$

Before moving to several computations on those observables, we state general regularity facts. We say that a function is $t$-local if $f(g, o)$ only depends on $(g, o)_t$.

Recall that $\vartheta_1 = L/d$.

Lemma 12.1 The function $f_{\phi, \psi, t}$ is $t$-local and satisfies

$$|f_{\phi, \psi, t}(g, o)| \leq |\phi|_\infty |\psi|_\infty |P_g(o, t)|_1 \vartheta_1^t$$ \hspace{1cm} (12.4)

The function $F_{\psi, t}$ is $(t + 1)$-local and satisfies

$$|F_{\mu, \psi, t}(g, o)| \leq |\psi|_\infty |P_g(o, t + 1)|((\vartheta_1^t + \vartheta_1^{t+1})/|\mu|).$$ \hspace{1cm} (12.5)

Proof The locality property is obvious from the definition, while for the bound it suffices to write

$$|f_{\phi, \psi, t}(g, o)| \leq \left(\frac{n}{d}\right)^t |\phi|_\infty \sum_{P_g(o, t)} \left(\frac{L}{n}\right)^t |\psi|_\infty$$

$$\leq \vartheta_1^t |\phi|_\infty |\psi|_\infty |P_g(o, t)|.$$
It is the same thing for $F_{\psi,t}$.

The following crucial theorem gathers all the computations linked with expectations or variances of those functionals when specialized on a tree $(T_x, x)$ with the distribution described before.

**Theorem 12.2** Let $\psi$ be any vector in $\mathbb{R}^n$ and $t$ be an integer. For any $i, j \in [r]$, the following identities are true.

\[
E[f_{\psi,\varphi_i,t}(T, x)] = \psi(x)\varphi_j(x)\mu_j^t \quad (12.6)
\]

\[
E[f_{\psi,\varphi_i,t}(T, x)f_{\psi,\varphi_j,t}(T, x)] = \mu_i^t\mu_j^t\psi(x)^2\sum_{s=0}^t \frac{Q^s \varphi_{i,j}(x)}{(\mu_i\mu_j)^s}. \quad (12.7)
\]

\[
E[F_{\mu_i,\varphi_i,t}(T, x)^2] = \frac{Q^t \varphi_{i,i}(x)}{dt}. \quad (12.8)
\]

The proof consists in using the eigenvector equation to identify specific martingales and take advantage of their properties to compute those expectations and variances. It is postponed to Sect. 14.

### 12.2 Functionals on Graphs: Concentration

This section describes concentration of sum functionals on the graph $G$, having the form $\sum_{o \in [n]} f(G, o)$ where $f: \mathcal{G} \to \mathbb{R}$ is any measurable function. The tools and spirit of this section are identical to [14, Section 9], but slightly adapted to our needs.

The first proposition deeply exploits the fact that $G$ is in fact a function of independent random variables $((M_{y,x}, M_{y,x}))_{y \geq x}$. A generalized Efron-Stein inequality will be very useful here.

**Proposition 12.3** (Moment inequality for graph functionals) Let $f, \bar{f}: \mathcal{G} \to \mathbb{R}$ be two $t$-local functions such that $|f(g, o)| \leq \bar{f}(g, o)$ and $\bar{f}$ is non-decreasing by the addition of edges. Then, for some universal constant $c > 0$, for all $p \geq 2$,

\[
\left( E\left[\left|\sum_{o \in [n]} f(G, o) - E\sum_{o \in [n]} f(G, o)\right|^p\right]\right)^{1/p} \leq c\sqrt{n}p^{3/2}D^t \left( E\left[\max_{x \in [n]} \bar{f}(G, x)^{2p}\right]\right)^{1/(2p)},
\]

and

\[
\left( E\left[\left|\sum_{o \in [n]} f(G, o) - E\sum_{o \in [n]} f(G, o)\right|^p\right]\right)^{1/p} \leq c\sqrt{n}p (p + \ln(n)) D^t \left( E\left[\frac{1}{n}\sum_{x \in [n]} \bar{f}(G, x)^{2p}\right]\right)^{1/(2p)}.
\]
Proof We define $E_x$ as the set of edges of the form $(x, y)$ or $(y, x)$ with $x \leq y$. From our assumptions on $M$, the variables $(E_x)_{x \in [n]}$ are independent. Moreover, there is a measurable function $F$ such that

$$\sum_{o \in [n]} f(G, o) = F(E_1, \ldots, E_n).$$

Let us denote $Y = F(E_1, \ldots, E_n)$ this sum, and for any $x$ let us note $Y_x$ the same sum where $E_x$ has been emptied:

$$Y_x = F(E_1, \ldots, E_{x-1}, \emptyset, E_{x+1}, \ldots, E_n).$$

Equivalently, if $G_x$ indicates the graph $G$ where all the directed edges between $x$ and a larger or equal vertex have been deleted, we have

$$Y_x = \sum_{o \in [n]} f(G_x, o).$$

The moment inequality [17, Theorem 15.5] implies that there exists a universal constant $0 < c < 6$ such that for all $p \geq 2$,

$$\mathbb{E} \left[ \sum_{o \in [n]} f(G, o) - \mathbb{E} \sum_{o \in [n]} f(G, o) \right]^p \leq (c \sqrt{p})^p \mathbb{E} \left[ \sum_{x \in [n]} (Y - Y_x)^2 \right].$$

This is a generalization of Efron-Stein inequality (corresponding to $p = 2$).

Now, fix $o$ and $x$ in $[n]$. The function $f$ is $t$-local, hence $f(G, o) - f(G_x, o)$ is always zero, except possibly if $x$ is in $(G, o)_t$, or equivalently if $o$ is in $(G, x)_t$. As a consequence, we have

$$|Y - Y_x| \leq \sum_{o \in (G, t)} f(G, o) + f(G_x, o) \leq 2|G, o_t| \max_{o \in [n]} \bar{f}(G, o)$$

where in the last line we used the fact that $\gamma$ is non-decreasing by the addition of edges. Recall that Hölder inequality implies that for all $p \geq 2$,

$$\left( \sum_{i=1}^{n} u_i^2 \right)^{p/2} \leq n^{p/2 - 1} \left( \sum_{i=1}^{n} |u_i|^p \right).$$

Therefore,

$$\mathbb{E} \left[ \sum_{x \in [n]} (Y - Y_x)^2 \right]^{p/2} \leq n^{p/2 - 1} 2^p \mathbb{E} \left[ \sum_{x \in [n]} \max_{o \in [n]} \bar{f}(G, o)^p |(G, x)_t|^p \right].$$
\[ \leq n^{p/2} \sqrt{E[(G, x)_t]^2 p} E \left[ \max_{o \in \mathcal{O}} \tilde{f}(G, o)^2 p \right]. \]

where we have used Cauchy-Schwarz inequality at the second line. Finally, we use 11.8 and it concludes the proof of the first statement of the proposition.

For the second statement, we write instead:

\[
\sum_{x \in [n]} (Y - Y_x)^2 \leq 4 \sum_{x \in [n]} \left( \sum_{o \in (G, x)_t} \tilde{f}(G, o) \right)^2
\]

\[
\leq 4 \sum_{x \in [n]} \left| (G, x)_t \right| \sum_{o \in (G, x)_t} \tilde{f}(G, o)^2
\]

\[
= 4 \sum_{o \in [n]} \tilde{f}(G, o)^2 \sum_{x \in (G, o)_t} \left| (G, x)_t \right|
\]

\[
\leq 4 \max_{x \in [n]} \left| (G, x)_t \right|^2 \sum_{o \in [n]} \tilde{f}(G, o)^2.
\]

The rest of the proof follows exactly the same line. \(\square\)

The next immediate lemma is a comparison principle between the expectation of a graph functional on the random graph \(G\) and the same functional on the random tree.

**Lemma 12.4** Let \(0 \leq \kappa \leq 0.49\) and \(h\) an integer as in (11.10). Let \(f : \mathcal{G}_* \to \mathbb{R}\) be a \(h\)-local function. Then, some universal constant \(c > 0\), we have for all \(x \in [n],\)

\[
|E f(G, x) - E f(T, x)| \leq c \ln(n) n^{\kappa - 1/2} \sqrt{E[f(G, x)^2]} + E[f(T, x)^2]. \quad (12.9)
\]

**Proof** Let \(E(x)\) denote the event “the coupling between \((G, x)_h\) and \((T, x)_h\) fails”; as our functionals are \(h\)-local, we have \(f(G, x) = f(T, x)\) on \(E(x)\). Proposition 11.4 implies that \(P(E(x)) \leq c (\ln n)^2 n^{2 \kappa - 1}\). Consequently, by the Cauchy-Schwarz inequality,

\[
|E f(G, x) - E f(T, x)| \leq E[|f(G, x) - f(T, x)|1_{E(x)}]
\]

\[
\leq \sqrt{P(E(x))} \left( \sqrt{E[f(G, x)^2]} + \sqrt{E[f(T, x)^2]} \right)
\]

\[
\leq \sqrt{c \ln(n) n^{\kappa + 1/2}} \left( \sqrt{E[f(G, x)^2]} + \sqrt{E[f(T, x)^2]} \right),
\]

which is bounded by the RHS in the claim, upon adjusting the constant. \(\square\)

Proposition 12.3 and Lemma 12.4 can be combined to derive general deviation inequalities for graph functionals. For simplicity in the next theorem, we consider the specific case of majorizing functions \(\tilde{f}(g, o)\) that we will encounter in the sequel.
**Theorem 12.5** Let $0 \leq \kappa \leq 0.49$ and $h$ an integer as in (11.10). Let $f : G^* \to \mathbb{R}$ be a $h$-local function such that $|f(g, o)| \leq \alpha |(g, o)_h|^{\beta}$ for some $\alpha, \beta > 0$. Then, for some universal constant $c > 0$, for any $s \geq 1$, with probability greater than $1 - n^{-s}$, we have

$$\left| \sum_{x \in [n]} f(G, x) - E \sum_{x \in [n]} f(T, x) \right| \leq ce^{\beta \alpha s^{3/2+\beta}} \ln(n)^{3/2+\beta} n^{\kappa(1+\beta)+1/2}.$$ 

**Proof** We set $\tilde{f}(g, o) = \alpha |(g, o)_h|^{\beta}$. By (11.8), we have for all $p \geq 1$ and $x \in [n]$, for some universal constant $c > 0$,

$$\mathbb{E}[\tilde{f}(G, x)^{2p}] \leq \alpha^{2p} (cp)^{2p\beta} D^{2p\beta h}.$$ 

By the Chebyshev inequality and the second claim of Proposition 12.3, adjusting the constant $c > 0$, we have for all $t > 0$,

$$\left| \sum_{x \in [n]} f(G, x) - E \sum_{x \in [n]} f(G, x) \right| \geq c\alpha sn^{\kappa(1+\beta)+1/2},$$ 

with probability at most

$$\left( \frac{p^{1/2+\beta}(p \vee \ln(n))}{t} \right)^p.$$ 

We take $t = (es \ln(n))^{3/2+\beta}$ with $s \geq 1$ and $p = t^{1/(3/2+\beta)} / e \geq \ln(n)$. We obtain a bound with probability at least $1 - n^{-s}$. Then we use Lemma 12.4 and up to adjusting the universal constant, we obtain the desired bound. $\Box$

### 13 Near Eigenvectors: Proof of Theorem 9.2

In this Section we prove Theorem 9.2, using the tools introduced earlier. For some $0 < \kappa < 1$ which will be fixed at the end, we set

$$\ell = \lfloor \frac{\kappa}{2} \log_D(n) \rfloor.$$ 

Here is the route taken: first, we prove different propositions related with precise bounds for the entries of the matrices $U$, $V$ or $\Phi^* A^\ell \Phi$. Often, the error terms look like

$$c_0(b \ln n)^{c_1} n^{2\kappa - 1/2} \delta^t,$$

or small variants. For a good choice of $\kappa$, this gives the requested bounds in Theorem 9.2.
For functionals such as \( \langle \varphi_i, A^t \varphi_j \rangle \), the plan is simple: we justify why those functionals can be well-approximated by the identities of Theorem 12.2 thanks to the deviation inequality Theorem 12.5.

Bounding \( \| A^t \mathcal{P}_{H^\perp} \| \) is however much more difficult and will be done through a tangle-free decomposition, in Sect. 13.3. Performing the expected high-trace method requires some care and we postponed this part to Sect. 15.

13.1 Entry-Wise Bounds for Theorem 9.2

Proposition 13.1 Assume \( 0 \leq \kappa \leq 0.33 \). There is a universal constant \( c > 0 \) such that, with probability greater than \( 1 - cn^{3\kappa - 1} \), for any \( i, j \in [r_0] \) and \( t \leq 3\ell \), the following holds:

\[
\left| \langle \varphi_i, A^t \varphi_j \rangle - \mu^t_i \delta_{i,j} \right| \leq cb^2 (\ln n)^{5/2} n^{3\kappa - 1/2} \delta_1^t.
\] (13.1)

Proof Fix \( i, j \in [r_0] \) and \( t \leq 3\ell \). Using the notation already introduced in (12.1), we define a function \( f \) by

\[
f(g, o) = 1_{(g, o)_t} \text{ is tangle free } f_{\varphi_i, \varphi_j, t}(g, o).
\]

This function is clearly \( t \)-local and from 12.4, 4.3, and 11.1,

\[
|f(g, o)| \leq |\varphi_i|_\infty |\varphi_j|_\infty \delta_1^t |\mathcal{P}_g(o, t)|1_{(g, o)_t} \text{ is tangle free} \leq \frac{2b^2}{n} \delta_1^t |(g, o)_t|.
\]

Moreover, on the event that \( G \) is \( t \)-tangle free, from (12.2) we have

\[
\langle \varphi_i, A^t \varphi_j \rangle = \sum_{x \in [n]} f(G, x)
\]

We now apply the concentration result in Theorem 12.5 to the function \( f \) with \( s = 4 \), \( \alpha = 2b^2 \delta_1^t / n \), \( \beta = 1 \) and \( \kappa' = 3\kappa / 2 \). The error bound in Theorem 12.5 is thus, for some new constant \( c > 0 \),

\[
cb^2 (\ln n)^{5/2} n^{3\kappa - 1/2} \delta_1^t.
\]

Moreover, as computed in Theorem 12.2 - Eq. (12.6), we have

\[
\mathbf{E} \sum_{x \in [n]} f(T, x) = \mathbf{E} \sum_{x \in [n]} f_{\varphi_i, \varphi_j, t}(T_x, x) = \mu^t_i \delta_{i,j}.
\]

Combined with Lemma 11.3 to control the probability that the graph \( 3\ell \)-tangle free and \( \ell \leq \ln(n) \) for the union bound, this concludes the proof. \( \square \)
**Proposition 13.2** Assume \( 0 \leq \kappa \leq 0.49 \). There is a universal constant \( c > 0 \) such that, with probability greater than \( 1 - cn^{2\kappa - 1} \), for any \( i, j \in [r_0] \) and \( t \leq \ell \), the following holds:

\[
\left| \langle A^t \varphi_i, A^t \varphi_j \rangle - \mu_{t,i}^l \mu_{j,t}^l \Gamma_{i,j}^{(t)} \right| \leq cb^2 (\ln n)^{7/2} n^{3\kappa/2 - 1/2} \frac{\delta_1^{2t}}{n^2} \tag{13.2}
\]

\[
\left| \langle (A^*)^t \varphi_i, (A^*)^t \varphi_j \rangle - \mu_{t,i}^l \mu_{j,t}^l \Gamma_{i,j}^{(t)} \right| \leq cb^2 (\ln n)^{7/2} n^{3\kappa/2 - 1/2} \frac{\delta_1^{2t}}{n^2} \tag{13.3}
\]

**Proof** Since \( A^* \) and \( A \) are identical in distribution, one only has to prove the first inequality. The proof is the same as for Proposition 13.1. Fix \( i, j \in [r_0] \) and \( t \leq \ell \).

The right function here is \( f \) defined by

\[
f(g, o) = 1_{(g,o)_t} \text{ is tangle free } f_1, \varphi_i, t(g, o) f_1, \varphi_j, t(g, o).
\]

This function is clearly \( t \)-local and from 12.4, 4.3, and 11.1,

\[
|f(g, o)| \leq |\varphi_i|_{\infty} |\varphi_j|_{\infty} \delta_1^{2t} |P_g(o, t)|^2 1_{(g, o)_t} \text{ is tangle free} \leq \frac{4b^2}{n} \delta_1^{2t} |(g, o)_t|^2.
\]

We observe that if the graph \( G \) is \( t \)-tangle free, we have

\[
\langle A^t \varphi_i, A^t \varphi_j \rangle = \sum_{x \in [n]} f(G, x).
\]

We now apply Theorem 12.5 with \( s = 4, \alpha = 4b^2 \delta_1^{2t}/n, \beta = 2 \) and \( \kappa' = \kappa/2 \). The error bound in Theorem 12.5 is thus, for some new \( c > 0 \),

\[
\frac{c}{b^2} (\ln n)^{7/2} n^{3\kappa/2 - 1/2} \frac{\delta_1^{2t}}{n^2}.
\]

Moreover, as computed in Theorem 12.2 - Eq. (12.7), we have

\[
\mathbf{E} \left[ \sum_{x \in [n]} f(T, x) \right] = \mathbf{E} \left[ \sum_{x \in [n]} f_1, \varphi_i, t(T_x, x) f_1, \varphi_j, t(T_x, x) \right] = \mu_{t,i}^l \mu_{j,t}^l \Gamma_{i,j}^{(t)}.
\]

Combined with Lemma 11.3, this concludes the proof. \( \square \)

### 13.2 Control over the Growth of a Process

In this section, we establish the following proposition. In words, it asserts that if \( \langle (A^*)^t \varphi_i, w \rangle = 0 \) then \( \langle (A^*)^t \varphi_i, w \rangle \) is quite small for all \( t \leq \ell \).
Proposition 13.3 Assume $0 \leq \kappa \leq 0.33$. There exists a universal constant $c > 0$ such that, with probability greater than $1 - cn^{2\kappa-1}$, one has for any $t \leq \ell$, for any $w \in H^\perp$ with $|w| = 1$ and for any $i \in [r_0]$ the following bound:

$$|\langle (A^*)^t \varphi_i, w \rangle| \leq c \ell L^2 \delta^t.$$ 

Proof We follow the usual strategy. First, we note that by the mere definition of $H$, when $w \in H^\perp$ we have $\langle (A^*)^\ell \varphi_i, w \rangle = 0$. Consequently,

$$\mu_i^{-t} \langle (A^*)^t \varphi_i, w \rangle = \mu_i^{-t} \langle (A^*)^t \varphi_i, w \rangle - \mu_i^{-\ell} \langle (A^*)^\ell \varphi_i, w \rangle$$

and from a telescopic sum we get

$$|\mu_i^{-t} \langle (A^*)^t \varphi_i, w \rangle| = \left| \sum_{k=1}^{\ell-1} \mu_i^{-k} \langle (A^*)^k \varphi_i, w \rangle - \mu_i^{-k} \langle (A^*)^{k+1} \varphi_i, w \rangle \right|$$

$$\leq \sum_{k=t}^{\ell-1} |\mu_i|^{-k} \langle (A^*)^k \varphi_i - \frac{1}{\mu_i} (A^*)^{k+1} \varphi_i \rangle$$  \hspace{1cm} (13.4)

where in the last line we used the Cauchy-Schwarz inequality and $|w| = 1$. Let us fix $1 \leq k \leq \ell - 1$. We have

$$\left| (A^*)^k \varphi_i - \frac{1}{\mu_i} (A^*)^{k+1} \varphi_i \right|^2 = \sum_{x \in [n]} F_{\mu_i, \varphi_i, k}(G, x)^2$$

where $F_{\mu_i, \varphi_i, k}$ was defined in (12.3). From Theorem 12.2 - Eq. (12.8), we have for any $x \in [n]$

$$E[F_{\mu_i, \varphi_i, k}(T, x)^2] = \frac{Q^{k \varphi^{i,i}(x)}}{d^k}.$$ 

Therefore, summing over $x$ gives

$$E \left[ \sum_{x \in [n]} F_{\mu_i, \varphi_i, k}(T, x)^2 \right] = \frac{\langle 1, Q^{k \varphi^{i,i}} \rangle}{d^k}.$$ 

We are going to use the concentration bound from Theorem 12.5. We define the function $f$ by

$$f(g, o) = 1_{(g, o)_{k+1}} \text{ is tangle free } F_{\mu_i, \varphi_i, k}(g, o)^2.$$
This function is \((k+1)\)-local and from 12.5, 4.3, and 11.1, we have

\[
    f(g, o) \leq |\varphi_i|_\infty^2 |\mathcal{P}_g(o, k+1)|^2 (\delta_1^k + \delta_1^{k+1}/|\mu_i|)^2 1_{(g, o)_{k+1}} \text{ is tangle free} \quad (13.5)
\]

\[
    \leq \frac{8b^2}{n} \delta_1^{2k} |(g, o)_{k+1}|^2, \quad (13.6)
\]

where we have used \(|\mu_i| \geq \delta_1\). Then, applying Lemma 11.3 and Theorem 12.5 (with \(s = 4, \alpha = 8b^2 \delta_1^{2k}/n\) and \(\beta = 2\)) yields, with probability at least \(1 - cn^{2^k - 1}\), for all \(k \leq \ell - 1\) and \(i \in [r_0]\),

\[
    \left| (A^*)^k \varphi_i - \frac{1}{\mu_i} (A^*)^{k+1} \varphi_i \right|^2 - \frac{\langle 1, Q_k^i \varphi_i^j \rangle}{d^k} \leq c^2 b^2 (\ln n)^{7/2} n^{3^k/2 - 1/2} \delta_1^{2k}.
\]

We now sum all \(k\) between \(t\) and \(\ell - 1\) in 13.4. We find

\[
    |\mu_i^{-t} \langle (A^*)^t \varphi_i, w \rangle | \leq \sum_{k=t}^{\ell-1} \left( \frac{\langle 1, Q_k^i \varphi_i^j \rangle}{(\mu_i^2 d)^k} + cb(\ln n)^{7/4} n^3^{k/4 - 1/4} \sum_{k=t}^{\ell-1} \delta_1^k \right) \frac{K^2 \delta_1^k}{|\mu_i|^k},
\]

where we used \(\sqrt{u + v} \leq \sqrt{u} + \sqrt{v}\). Then, we use 10.5 which tells that \(\langle 1, Q_k^i \varphi_i^j \rangle \leq K^2 \delta_1^k\). We thus get from 13.4 the inequality

\[
    |\langle (A^*)^t \varphi_i, w \rangle | \leq |\mu_i|^t \left( K + cb(\ln n)^{7/4} n^{3^k/4 - 1/4} \sum_{k=t}^{\ell-1} \delta_1^k \right) \frac{K^2 \delta_1^k}{|\mu_i|^k}
\]

\[
    \leq (\ell - t)(K + cb(\ln n)^{7/4} n^{3^k/4 - 1/4}) \delta_1^t,
\]

where we used the fact that \(|\mu_i| > \delta_1\). We finally use 10.2, that is \(K \leq L^2/\mu_1^2 = L^2\) since we supposed that \(\mu_1 = 1\). \(\square\)

13.3 Norm of the Matrix Restricted to \(H^k\)

Our goal in this section is to prove the inequality 9.11 on \(\|A^{(k)}\|\). We first describe the tangle-free decomposition introduced in [14, 43]. The main technical estimate will be postponed to Sect. 15.

We notice that if the graph \(G\) is \(\ell\)-tangle free then \(A^\ell = A^{(\ell)}\) where

\[
    A^{\ell}_{x,y} = \left( \frac{n}{d} \right)^\ell \sum_{F^\ell_{x,y}} \prod_{t=1}^{\ell} P_{x_{t-1} x_t} M_{x_{t-1} x_t},
\]

and the sum runs over the set \(F^\ell_{x,y}\) of all paths \((x_0, \ldots, x_\ell)\) such that \(x_0 = x, x_\ell = y\) and the graph of the path is tangle-free — we recall that tangle-free means that there are no more than one cycle, see the definitions in Sect. 11.1 on page 50. More generally,
$F^t$ denotes the set of all tangle-free paths of length $t$, whatever their endpoints. We also define the matrices $M$ and $A^{(\ell)}$ by $A^{(0)} = A^{(0)} = I_n$, and

$$M_{x,y} = M_{x,y} - \frac{d}{n},$$

$$A^{(\ell)}_{x,y} = \left(\frac{n}{d}\right)^{\ell} \sum_{F^t_{x,y}} P_{x_{t-1}x_t} M_{x_{t-1}x_t}, \quad (13.7)$$

We use the convention that the product over an emptyset is 1. Then we may write for any $a, b \in \mathbb{R}^\ell$,

$$\prod_{t=1}^\ell a_t = \prod_{t=1}^\ell b_t + \sum_{k=1}^{\ell-1} \left(\prod_{t=1}^k (a_k - b_k)\right) \left(\prod_{t=k+1}^\ell a_t\right). \quad (13.8)$$

We thus get

$$A^{(\ell)}_{x,y} = A^{(\ell)}_{x,y} + \sum_{k=1}^\ell \left(\frac{n}{d}\right)^{\ell} \sum_{F^t_{x,y}} P_{x_{t-1}x_t} M_{x_{t-1}x_t} \left(\frac{d}{n} P_{x_{k-1}x_k}\right) \prod_{t=k+1}^\ell P_{x_{t-1}x_t} M_{x_{t-1}x_t}. \quad \text{(13.8)}$$

This can then be rewritten as the following identity in $\mathcal{M}_n(\mathbb{R})$:

$$A^{(\ell)} = A^{(\ell)} + \sum_{k=1}^\ell A^{(k-1)} PA^{(\ell-k)} - \sum_{k=1}^\ell R^{(\ell)}_k,$$

where

$$R^{(\ell)}_k = \left(\frac{n}{d}\right)^{\ell-1} \sum_{(x_0, \ldots, x_\ell) \in F^k_{x,y}} \prod_{t=1}^{k-1} P_{x_{t-1}x_t} M_{x_{t-1}x_t} P_{x_{k-1}x_k} \prod_{t=k+1}^\ell P_{x_{t-1}x_t} M_{x_{t-1}x_t},$$

where the sum is over all ‘paths’ $(x_0, \ldots, x_\ell)$ such that $(x_0, \ldots, x_{k-1}) \in F^{k-1}$, $(x_k, \ldots, x_\ell) \in F^{\ell-k}$ but $(x_0, \ldots, x_\ell)$ is not in $F^\ell$.

We now use the spectral decomposition $P = \mu_1 \varphi_1 \varphi_1^* + \cdots + \mu_n \varphi_n \varphi_n^*$. For any unit vector $w$, we have

$$A^{(k-1)} PA^{(\ell-k)} w = \sum_{j=1}^n \mu_j \varphi_j \langle \varphi_j, A^{(\ell-k)} w \rangle.$$

Hence, from the orthogonality of the $\varphi_j$’s,

$$|A^{(k-1)} PA^{(\ell-k)} w| \leq \left| A^{(k-1)} \right| \left( \sum_{j=1}^n \left| \mu_j \varphi_j \langle \varphi_j, A^{(\ell-k)} w \rangle \right| \right).$$
\[ = \|A^{(k-1)}\| \sqrt{\sum_{j=1}^{n} \mu_j^2 \langle \varphi_j, A^{(\ell-k)}w \rangle^2}. \] (13.9)

From Proposition 13.3, with probability at least \(1 - cn^{2\kappa-1}\), the following holds for any \(t \leq \ell\) and \(i \in [r_0]\) and \(w \in H^\perp\):

\[ |\langle \varphi_i, A^t w \rangle| \leq cb^4 \ell^t. \]

From Proposition 13.2, with probability at least \(1 - cn^{2\kappa-1}\), for all \(i \in [n]\setminus[r_0]\) and \(w \in H^\perp\):

\[ |\langle \varphi_i, At w \rangle| \leq \mu_2 t_i / \Gamma_1(t_i + \vartheta \ell - k) \leq \mu_2 t_i / \Gamma_1(t_i + \vartheta \ell - k + \ell). \]

However, from Eqs. (10.7–10.12), we have \(\mu_2 t_i / \Gamma_1(t_i + \vartheta \ell - k + \ell) \leq L^4(t + 1)\vartheta^{2t}\). As a consequence, for some universal constant \(c > 0\), for all \(i \in [n]\setminus[r_0]\) and \(t \leq \ell\),

\[ |\langle \varphi_i, A^t w \rangle| \leq |w| |(A^*)^t \varphi_i| \leq cL^2 \sqrt{\ell} \vartheta^{\ell - k}. \]

On the union of the two events events and \(G \ell\)-tangle free, the whole square root in 13.9 is bounded by, for all \(w \in H^\perp\):

\[ \sqrt{\sum_{j=1}^{r} \mu_j^2 \langle \varphi_j, A^{(\ell-k)}w \rangle^2} \leq \sum_{j=1}^{n} \mu_j^2 = cL^2 \ell \vartheta^{\ell - k} \leq \frac{cL^2 \ell \vartheta^{\ell - k}}{2}. \]

where we used that \(\mu_1 = 1\). We get the following lemma.

**Lemma 13.4** With probability at least \(1 - cn^{2\kappa-1}\), one has

\[ \|A^\ell P_{H^\perp}\| \leq \|A^\ell\| + cL^2 \ell \vartheta^{\ell - k} \sum_{k=1}^{\ell} \|A^{(k-1)}\| \vartheta^{\ell - k} + \sum_{k=1}^{\ell} \left\| R_k^{(\ell)} \right\|. \]

We now need bounds on \(\|A^{(k-1)}\|\) and \(\|R_k^{(\ell)}\|\).

**Proposition 13.5** There exists a universal constant \(c > 0\) such that if \(n \geq cK^{42}\), with probability at least \(1 - 1/\sqrt{n}\), the following holds for any \(k \in [\ell]\):

\[ \|A^{(k)}\| \leq c \ln(n)^{10} K^4 \vartheta^k. \] (13.10)

\[ \|R_k^{(\ell)}\| \leq \frac{cd}{n} \ln(n)^{23} L^\ell. \] (13.11)

The proof of this proposition relies on a high-trace method. It is postponed to Sect. 15. As a corollary, we obtain the following proposition.
Proposition 13.6 Assume $0 \leq \kappa \leq 1/4$. There exists a universal constant $c > 0$ such that, with probability greater than $1 - c/\sqrt{n}$, one has

$$\| A^\ell P_{H^\perp} \| \leq c L^{10} r \ln(n)^{12} \vartheta^\ell.$$ 

Proof By Lemma 13.4 and Proposition 13.5, with probability at least $1 - c/\sqrt{n}$, we have

$$\| A^\ell P_{H^\perp} \| \leq \| A^{(\ell)} \| + c L^2 \ell r \sum_{k=1}^\ell \| A^{(k-1)} \| \vartheta^{\ell-k} + \sum_{k=1}^\ell \| R^{(\ell)}_k \|$$

$$\leq C_0 \ln(n)^{10} \sum_{k=0}^\ell \vartheta^\ell + \frac{4d \ell \ln(n)^{24} L^\ell}{n}$$

$$\leq C_0 \ln(n)^{10} \ell \vartheta^\ell + c \ell d \ell \ln(n)^{24} \vartheta^1 \frac{d \ell}{n}$$

$$\leq C \ln(n)^{12} \vartheta^\ell$$

with $C_0 := c'L^{10} \ell r$, for some universal constant $c'$, in the second line (recall that $K \leq L$ from (10.2) and $\mu_1 = 1$). To get the last line, we have observed $d \ell^2 + 1/n \leq \vartheta^\ell n^{\kappa/2 + 1/8 - 1}$ where we have used $d \leq n^{1/8}$ from 9.2 and $\kappa \leq 1/4$. The constant $C$ in this last line is taken to be $C = c L^{10} r$ with $c$ some absolute constant.

13.4 Proof of Theorem 9.2

We gather the events and bounds from the last propositions, working out the error terms and presenting them in a way which keeps track of dependencies with the parameters.

As in 9.1, we define the vector space $H' = \text{vect}(u_1, \ldots, u_{r_0}) = \text{im}(U)$. Since $A$ and $A^*$ have the same distribution, Proposition 13.6 holds for $(A^*)^\ell P_{H'^\perp}$ in place of $A^\ell P_{H^\perp}$. Then, we set $\kappa = 1/4$ so that $\ell = [(\kappa/2) \log_2(n)]$ is as in Theorem 9.2. We apply Propositions 13.1 and 13.2 and Proposition 13.6 for $A^\ell P_{H^\perp}$ and $(A^*)^\ell P_{H'^\perp}$ has probability greater than $1 - c n^{3\kappa-1} = 1 - c n^{-\kappa}$.

On this good event, we should check that the error terms 9.6–9.11 are as claimed in the statement of Theorem 9.2. The claims 9.11–9.12 are contained in Proposition 13.6. The proofs of (9.6)-(9.10) are essentially the same. They all rely on the basic inequality for $T \in \mathcal{M}_{r_0, r_0}(\mathbb{R})$

$$\| T \| \leq r_0 \max_{i,j} |T_{i,j}|$$

and the use of the events in Propositions 13.1 and 13.2 to control the individual entries of the matrices.
The \((i, j)\)-entry of the matrix \(V^* A^\ell U = D^{-\ell} \Phi^* A^{3\ell} \Phi D^{-\ell}\) is precisely \((\mu_i \mu_j)^{-\ell} \langle \varphi_i, A^{3\ell} \varphi_j \rangle\). Hence on the event of Proposition 13.1, we find

\[
\| D^{-\ell} \Phi^* A^{3\ell} \Phi D^{-\ell} - \Sigma^\ell \| \leq cr_0 b^2 \ln(n)^{5/2} n^{3\kappa - 1/2} \frac{\delta^{3\ell}}{|\mu_i|^{\ell} |\mu_j|^{\ell}}.
\]

Our choice of \(\kappa = 1/10\) implies that \(3\kappa - 1/2 = -\kappa\). We use that \(\delta \leq \tau_0 |\mu_i|\) for all \(i \in [r_0]\). Adjusting the constant \(c > 0\), this gives the requested bound in (9.6).

We now prove 9.7. The \((i, j)\) entry of \(U^* V\) is equal to

\[
(\mu_i \mu_j)^{-\ell} \langle A^\ell \varphi_i, (A^*)^\ell \varphi_j \rangle = (\mu_i \mu_j)^{-\ell} \langle A^{2\ell} \varphi_i, \varphi_j \rangle.
\]

Hence, on the event of Proposition 13.1, we find

\[
\| U^* V - I_{r_0} \| \leq cr_0 b^2 \ln(n)^{5/2} n^{3\kappa - 1/2} \tau_0^{2\ell}.
\]

This implies 9.7. The bound 9.8 is proven similarly using the event of Proposition 13.1 for \(\kappa\) equal to \(\kappa/3\). It gives

\[
\| \Phi^* V - I_{r_0} \| \leq cr_0 b^2 \ln(n)^{5/2} n^{\kappa - 1/2} \tau_0^{\ell}.
\]

It remains to use 9.3 and our choice of \(\kappa\).

We now prove 9.9–9.10. The \((i, j)\) entry of \(U^* U\) is equal to

\[
(\mu_i \mu_j)^{-\ell} \langle A^\ell \varphi_i, A^\ell \varphi_j \rangle.
\]

On the event of Proposition 13.2, we find

\[
\| U^* U - \Gamma^{(\ell)} \| \leq cr_0 b^2 \ln(n)^{7/2} n^{2\kappa - 1/2} \tau_0^{2\ell}.
\]

This implies 9.9–9.10 and concludes the proof of Theorem 9.2.

14 Eigenwaves on Galton-Watson Trees: Proof of Theorem 12.2

This section carries out the details of the expectation and variance computation in Theorem 12.2. Let us recall the notation, and especially the definition 12.1 of the functionals: if \(\psi, \phi\) are two vectors in \(\mathbb{R}^n\) and \(t\) is an integer, then

\[
\sum_{\mathcal{D}(o, t)} \Phi(t(o)) \sum_{\mathcal{P}_x(o, t)} P_{t(o), t(x_1)} \cdots P_{t(x_{t-1}), t(x_t)} \times \psi(t(x_t)).
\]
14.1 An Elementary Computation on Poisson Sums

Let $N$ be a $\text{Poi}(d)$ random variable, and let $(X_i), (Y_i)$ two i.i.d. sequences of random variables, both being independent from $N$; we suppose that $X_i$ is independent of $Y_j$ for $i \neq j$, but there might be a nontrivial dependence between $X_i$ and $Y_i$. Let us note

$$A = \sum_{i=1}^{N} X_i \quad B = \sum_{i=1}^{N} Y_i.$$ 

The following (classical) identity will be crucial in the next sections. For convenience, we provide a proof.

$$\text{Cov}(A, B) = d \mathbb{E}[XY]. \quad (14.1)$$

**Proof of 14.1** Primary computations show that $\mathbb{E}[A] = d \mathbb{E}[X]$ and $\mathbb{E}[B] = d \mathbb{E}[Y]$, hence

$$\text{Cov}(A, B) = \mathbb{E}[AB] - \mathbb{E}[A] \mathbb{E}[B] = \mathbb{E}[AB] - d^2 \mathbb{E}[X] \mathbb{E}[Y].$$

The first term $\mathbb{E}[AB]$ is thus equal to

$$\mathbb{E}[AB] = \mathbb{E} \left[ \sum_{i=1}^{N} \sum_{j=1}^{N} X_i Y_j \right]$$

$$= \sum_{k=0}^{\infty} \frac{e^{-d} d^k}{k!} \mathbb{E} \left[ \sum_{i=1}^{k} \sum_{j=1}^{k} X_i Y_j \right]$$

$$= \sum_{k=0}^{\infty} \frac{e^{-d} d^k}{k!} (k \mathbb{E}[XY] + k(k - 1) \mathbb{E}[X] \mathbb{E}[Y])$$

$$= \mathbb{E}[XY] \mathbb{E}[N] + \mathbb{E}[X] \mathbb{E}[Y] \mathbb{E}[N(N - 1)].$$

We have $\mathbb{E}[N(N - 1)] = d^2$, hence 14.1 holds true. \qed

Identity 14.1 will be used many times in the following context. Fix one vertex $x \in [n]$ and suppose that $X = P_{x,U} \varphi_i(U)$ and $Y = P_{x,U} \varphi_j(U)$ with $U \sim \text{Unif}[n]$. By the eigenvector equation $P \varphi_k = \mu_k \varphi_k$, we have $\mathbb{E}[X] = (\mu_i d \varphi_i(x))/n$ and $\mathbb{E}[Y] = (\mu_j d \varphi_j(x))/n$, hence in this case

$$\mathbb{E}[XY] = \frac{1}{n} \sum_{y \in [n]} P_{x,y}^2 \varphi_i(y) \varphi_j(y) = \frac{1}{n^2} (Q \varphi^{i,j})(x). \quad (14.2)$$

These identities will be used later in variance computations.

14.2 Proof of a Martingale Property

Let $x$ be a fixed element in $[n]$ and let $(T, x)$ be the random rooted marked tree described in Sect. 11.1 and let $\mathcal{F}_t$ be the sigma-algebra generated by $(T, x)_t$; from
now on we will use the filtration $\mathcal{F} = (\mathcal{F}_t)_{t \geq 0}$. The key observation for this whole section is that the process $t \mapsto \mu_k^{-t} f_{\phi, \varphi_k, t}(T, x)$ is indeed an $\mathcal{F}$-martingale.

**Lemma 14.1** Let $\phi$ be any vector and let $\varphi_i$ be an eigenvector of $P$ associated with the nonzero eigenvalue $\mu_i$. Then, the discrete-time stochastic process

$$Z_t := \frac{1}{\mu_i} f_{\phi, \varphi_i, t}(T, x)$$

is an $\mathcal{F}$-martingale.

From now on, the conditional expectation with respect to the sigma-algebra $\mathcal{F}_t$ will be noted $\mathbb{E}_t$ instead of $\mathbb{E}[\cdot | \mathcal{F}_t]$.

**Proof** It is clear that $f_{\phi, \varphi_i, t}(T, x) = \phi(x) f_{1, \varphi_i, t}(T, x)$, hence it is sufficient to prove the martingale property for $Z_t = \mu_i^{-t} f_{1, \varphi_i, t}(T, x)$.

Let us fix an integer $t$. Then, upon factorizing up to depth $t$ we have

$$Z_{t+1} - Z_t = \left( \frac{n}{d \mu_i} \right)^{t+1} \sum_{d^+(a, x) = t} \prod_{s=1}^t P_{t(x_s-1), t(x_s)} \left( \sum_{x_t \rightarrow y} P_{t(x_t), t(y)} \varphi_i(t(y)) - \frac{d \mu_i}{n} \varphi_i(t(x_t)) \right).$$

Let us note $\Delta_t = Z_{t+1} - Z_t$ the martingale increment. Then,

$$\mathbb{E}_t(\Delta_t) = \left( \frac{n}{d \mu_i} \right)^{t+1} \sum_{d^+(a, x) = t} \prod_{s=1}^t P_{t(x_s-1), t(x_s)} \mathbb{E}_t \left[ \sum_{x_t \rightarrow y} P_{t(x_t), t(y)} \varphi_i(t(y)) - \frac{d \mu_i}{n} \varphi_i(t(x_t)) \right].$$

Let $X_1, X_2, \ldots$ be i.i.d. random variables with the following distribution (conditionally on $\mathcal{F}_t$):

$$P_t(X = P_{t(x), z} \varphi_i(z)) = \frac{1}{n} \quad \text{for each } z \in [n].$$

In other words, conditionally on $\mathcal{F}_t$, the rv’s $X_s$ are i.i.d. samples with distribution $P_{t(x), U} \varphi_i(U)$ with $U \sim \text{Unif}[n]$, just as in the end of the preceding paragraph.

It is clear that for every children $y$ of $x_t$, the random variable $P_{t(x_t), t(y)} \varphi_i(t(y))$ has this distribution, and as already noted,

$$\mathbb{E}[X] = \frac{1}{n} \sum_{z \in [n]} P_{t(x), z} \varphi_i(z) = \frac{1}{n} (P \varphi_i)(t(x_t)) = \frac{\mu_i}{n} \varphi_i(t(x_t)).$$
The number \( N \) of children of \( x_t \) has a \( \text{Poi}(d) \) distribution, and is independent of \( F_t \), hence

\[
E_t \left[ \sum_{x_t \to y} P_{(x_t),t(y)} \varphi_i(t(y)) - \frac{d\mu_i}{n} \varphi_i(t(x_t)) \right] = E_t \left[ \sum_{s=1}^N X_s \right] - \frac{d\mu_i}{n} \varphi_i(t(x_t)) \]
\[
= dE[X] - \frac{d\mu_i}{n} \varphi_i(t(x_t)) \]
\[
= 0.
\]

We have \( E_t[\Delta_1] = 0 \) and the martingale property for \( Z_t \) is true.

\[\square\]

14.3 Proof of 12.6

The proof of these two identities is straightforward. Indeed, the martingale property for \( Z_t = \mu_t^{-1} f_{\psi,\varphi_i,t}(T, x) \) shows that

\[
E[f_{\psi,\varphi_i,t}(T, x)] = \frac{E[Z_t]}{\mu_i} = \frac{E[Z_0]}{\mu_i} = \psi(x) \varphi_i(x)
\]

which is exactly 12.6.

14.4 Proof of 12.7–12.8

We fix \( i, j \) in \([r]\) and \( x \) in \([n]\) for the rest of the proof. Clearly, it is enough to do the computations with \( \psi = 1 \). We set

\[
Z^i_t = \frac{f_{1,\varphi_i,t}(T, x)}{\mu_i}
\]

and

\[
\Delta_t = E_t[(Z^i_{t+1} - Z^i_t)(Z^j_{t+1} - Z^j_t)].
\]

The \( Z^i \) are martingales, hence

\[
E[Z^i_t Z^j_t] = E[\Delta_0 + \cdots + \Delta_{t-1}] = E[\Delta_0] + \cdots + E[\Delta_{t-1}].
\]

Our goal is to compute those \( \Delta_s \). First, we have

\[
\Delta_t = \left( \frac{n^2}{\mu_i \mu_j d^2} \right)^{t+1} \sum_{d^+(o,x) = t}^{t-1} \prod_{s=0}^{t-1} P_{(x_s),t(x_{s+1})} P_{(x'_s),t(x'_{s+1})} \times E(x_t, x'_t) \quad (14.3)
\]
where the sum runs over all the couples of paths of length \( t \) started at the root: \( o = x_0 \rightarrow x_1 \rightarrow \cdots x_t \) and \( x'_0 = o \rightarrow x'_1 \rightarrow \cdots \rightarrow x'_t \), and where \( E(x_t, x'_t) \) is given by

\[
E_t \left[ \sum_{x_t \rightarrow y} P_{t(x_t), t(y)} \varphi_j(t(y)) - \frac{d}{n} \mu_j \varphi_j(t(x_t)) \right] \times \left[ \sum_{x'_t \rightarrow y'} P_{t(x'_t), t(y')} \varphi_k(t(y')) - \frac{d}{n} \mu_k \varphi_k(t(x'_t)) \right].
\]

We have already computed those expectations in 14.1. More precisely, when \( x_t \neq x'_t \), the content of the two parentheses inside the expectation are totally independent and centered, hence the only contributions to (14.3) correspond to the summands where \( x_t = x'_t \). In this case, 14.1 and 14.2 yields

\[
E(x_t, x_t) = \frac{d}{n^2} Q^{i,j}(t(x_t)).
\]

We thus have

\[
\Delta_t = \frac{d}{n^2} \left( \frac{n^2}{\mu_i \mu_j d^2} \right)^{t+1} \sum_{d^+(o,x)=t} \prod_{s=0}^{t-1} P_{t(x_s), t(x_{s+1})}^2 Q^{i,j}(t(x_s)). \tag{14.4}
\]

Since our goal is to compute \( \mathbb{E}[\Delta_s] \) for any \( s \), we now apply repeated conditioning: \( \mathbb{E}[\Delta_t] = \mathbb{E}[\mathbb{E}_0[\mathbb{E}_1[\ldots \mathbb{E}_s[\Delta_s] \ldots ]]. \) By the computations done earlier, it is easy to see that

\[
\mathbb{E}[\Delta_s] = \left( \frac{d}{n^2} \right)^{s+1} \left( \frac{n^2}{\mu_i \mu_j d^2} \right)^{s+1} Q^{s,i,j}(t(x)) = \frac{Q^{s,i,j}(t(x))}{(\mu_i \mu_j d)^s}.
\]

This directly gives identity (12.8). Identity (12.7) also readily follows:

\[
\mathbb{E}[f_{\phi, i}, t(T, x)f_{\phi, j}, t(T, x)] = \mu_i^t \mu_j^t \phi(x)^2 \mathbb{E}[Z_i^t Z_j^t] = \mu_i^t \mu_j^t \phi(x)^2 \sum_{s=0}^{t} \frac{(Q^{s,i,j})(t(x))}{(\mu_i \mu_j d)^s}.
\]

As requested.

15 Proof of Proposition 13.5

In this section, we prove Proposition 13.5. The proof relies on the expected high trace method introduced in random matrix theory by Füredi and Komlós [31] and on techniques developed in [14] for sparse random matrices.
15.1 Norm of $A^{(k)}$

In this section, we prove the following lemma.

**Lemma 15.1** There exists a universal constant $c \geq 3$ such that for all integers $1 \leq k \leq \ln(n)$ and $n \geq c K^{42},$

$$\left( E \left\{ \| A^{(k)} \|^{2m} \right\} \right)^{\frac{1}{2m}} \leq \ln(n)^9 K^4 \vartheta^k,$$

where $m = \ln(n/K^6)/(12 \ln(\ln(n))).$

From Markov inequality, Lemma 15.1 implies Eq. (13.10). We start the proof of Lemma 15.1 by the norm identities

$$\| A^{(k)} \|^{2m} = \| A^{(k)} A^{(k)*} \| = \left\| \left( A^{(k)} A^{(k)*} \right)^m \right\|.$$

From the trace formula, we get

$$\| A^{(k)} \|^{2m} \leq \text{tr} \left\{ \left( A^{(k)} A^{(k)*} \right)^m \right\} = \sum_{(x_1, \ldots, x_{2m})} \prod_{t=1}^m (A^{(k)}_{x_2t-1, x_{2t}} (A^{(k)}_{x_{2t+1}, x_{2t}}),$$

where the product is over all $(x_1, \ldots, x_{2m})$ in $[n]^{2m}$ and we have set $x_{2m+1} = x_1.$

From the definition of $A^{(k)}$ in (13.7), taking expectation, we get

$$E \| A^{(k)} \|^{2m} \leq \left( \frac{n}{d} \right)^{2km} \sum_{\gamma} \prod_{i=1}^{2m} \prod_{t=1}^k P_{\gamma, i-1, \gamma, i, t \gamma, i-1, \gamma, i, t},$$

where the sum is over all $\gamma = (\gamma_1, \ldots, \gamma_{2m})$ with $\gamma_i = (\gamma_{i,0}, \ldots, \gamma_{i,k}) \in F^k$ and the boundary conditions: for all $i \in [m],$

$$\gamma_{2i,0} = \gamma_{2i+1,0} \quad \text{and} \quad \gamma_{2i-1,k} = \gamma_{2i,k}$$

with $\gamma_{2m+1} = \gamma_1.$

We associate to an element $\gamma$ as above, a directed graph $G_\gamma = (V_\gamma, E_\gamma)$ with vertices $V_\gamma = \{ \gamma_{i,t} : 1 \leq i \leq 2m, 0 \leq t \leq k \}$ and edge set $E_\gamma = \{ (\gamma_{i,t-1}, \gamma_{i,t}) : 1 \leq i \leq 2m, 1 \leq t \leq k \}.$ This graph may have loops (edges of $E_\gamma$ of the form $(x, x)$) and inverse edges (pair of edges $(x, y)$ and $(y, x)$ in $E_\gamma$). From the above boundary conditions, the graph $G_\gamma$ is simply connected. In particular, the genus of $G_\gamma$ is non-negative:

$$| E_\gamma | - | V_\gamma | + 1 \geq 0.$$
Each oriented edge \( e \in E_\gamma \) has a multiplicity \( m_e \) defined as the number of times it is visited by \( \gamma \):

\[
m_e = \sum_{(i,t)\in[2m] \times [k-1]} 1_{(\gamma_{i,t}, \gamma_{i,t+1})=e}.
\]

By construction,

\[
\sum_{e\in E_\gamma} m_e = 2km.
\] (15.3)

We may now estimate the expectation on the right-hand side of (15.1). Recall that the random variables \( M_{xy} = (M_{xy} - d/n) \), \( x, y \), are iid, centered, bounded by 1 and with variance \( (d/n) - (d/n)^2 \). It follows that for any \( p \geq 1 \), \( |E[M_{xy}^p]| \leq d/n \). We deduce that

\[
E \prod_{i=1}^{2m} \prod_{t=1}^k P_{\gamma_i,t-1\gamma_i,t}^{m_{\gamma_i,t-1\gamma_i,t}} M_{\gamma_i,t-1\gamma_i,t} = \prod_{e\in E_\gamma} P_e^{m_e} E[M_{1e}^{m_e}] \leq \left(\frac{d}{n}\right)^{|E_\gamma|} \prod_{e\in E_\gamma} |P_e|^{m_e}.
\]

Moreover, the above expectation is zero unless all edges have multiplicity at least 2. From 15.1, we thus obtain that

\[
E||A^{(k)}||^{2m} \leq \sum_{\gamma\in W_{k,m}} \left(\frac{n}{d}\right)^{2km-|E_\gamma|} \prod_{e\in E_\gamma} |P_e|^{m_e},
\] (15.4)

where \( W_{k,m} \) is the set of paths \( \gamma \) as above such that each edge of \( E_\gamma \) is visited at least twice.

We now organize the sum (15.4) in terms of the topological properties of the paths. We introduce the equivalence class in \( W_{k,m} \), we write \( \gamma \sim \gamma' \) if there exists a permutation \( \sigma \in S_n \) such that \( \gamma' = \sigma \circ \gamma \), where \( \sigma \) acts on \( \gamma \) by mapping \( \gamma_{i,t} \) to \( \sigma(\gamma_{i,t}) \). We denote by \( V_{k,m} \) the set of equivalence classes. Obviously, \( |V_\gamma| \) and \( |E_\gamma| \) are invariant in each equivalence class. For \( a, s \) integers, we denote by \( W_{k,m}(s, a) \) the equivalence classes such that \( |V_\gamma| = s \) and \( |E_\gamma| = a \). From (15.2), \( W_{k,m}(s, a) \) is empty unless \( a - s + 1 \geq 0 \). Our first lemma is a rough estimate on \( W_{k,m}(s, a) \).

**Lemma 15.2** Let \( a, s \geq 1 \) be integers such that \( a - s + 1 \geq 0 \). We have

\[
|W_{k,m}(s, a)| \leq (2km)^{6m(a-s+1)+2m}.
\]

**Proof** This lemma is contained in the proof of [14, Lemma 17]. We reproduce the proof for the reader convenience. Let \( \gamma = (\gamma_1, \ldots, \gamma_{2m}) \in W_{k,m} \). We order the set \( T = \{(i, t) : 1 \leq i \leq 2m, 0 \leq t \leq k-1\} \) with the lexicographic order. We think of \( T \) as time. For \( 0 \leq t \leq k-1 \) and \( i \) odd, we define \( e_{i,t} = (\gamma_{i,t}, \gamma_{i,t+1}) \), \( \gamma_{i,t} = \gamma_{i,t+1} \), while for \( i \) even, we set \( e_{i,t} = (\gamma_{i,k-t-1}, \gamma_{i,k-t}) \), \( \gamma_{i,t} = \gamma_{i,k-t-1} \) (in words: we reverse
γ_τ for even i). A vertex \( x \in V_γ \setminus \{γ_1, 1\} \) is visited for the first time at \( τ \in T \) if \( γ_τ = x \) and for all smaller \( σ \in T, γ_σ \neq x \).

We pick a distinguished path in each equivalence class by saying that \( γ \in W_{k,m} \) is canonical if \( V_γ = \{1, \ldots, |V_γ|\}, γ_1, 1 = 1 \) and vertices are first visited in order. There exactly one canonical path in each equivalence class. We thus aim for an upper bound on the number of canonical paths in \( W_{k,m} \) with \( |V_γ| = s \) and \( |E_γ| = a \) by designing an injective map (or encoding) on such canonical paths.

Our goal is to retrieve unambiguously the values of \( γ_τ, τ \in T \), from minimal information. For \( τ \in T \), we say that \( τ \) is a first time, if \( γ_τ \) has not been seen before. If \( τ \) is a first time the edge \( e_τ \) is called a tree edge. By construction, the set of tree edges is a sub-graph of \( G_γ \) with no weak cycle (without orientation) and vertex set \( V_γ \). We call the other edges of \( G_γ \) the excess edges. Any vertex different from 1 has its associated tree edge. It follows that the number of excess edges is

\[
g = a - s + 1.
\]

If \( e_τ \) is an excess edge, we say that \( τ \) is an important time. Other times are tree times (visit of a tree edge which has been seen before).

The set \( T_i = \{(i, t) : 0 \leq t \leq k - 1\} \) is composed by the successive repetitions of (i) a sequence of the tree times (possibly empty), (ii) a sequence of first times (possibly empty), (iii) an important time.

We build a first encoding of canonical paths. We mark the important times \( (i, t) \) by the vector \((γ_{i,t}, γ_{t-1})\), where \( τ \in T_i \cup \{(i, k)\} \) is the next time that \( e_τ \) will not be a tree edge (by convention \( τ = (i, k) \) if \( γ_i \) remains on the tree after \( (i, t) \)). We can reconstruct a canonical path \( γ \in W_{k,m} \), from the positions of the important times and their marks. Indeed, this follows from two observations (1) there is at most one path between two vertices in an oriented tree, and (2) if \( v \) vertices has been seen so far and \( τ \) is a first time then \( γ_τ = v + 1 \). It is our first encoding.

We refine this encoding by using the assumption that for each \( i, γ_i \) is tangle-free. We partition important times into three categories, short cycling, long cycling and superfluous times as follows. Assume that \( γ_i \) contains a cycle. Consider the smallest time \( (i, t_1) \) such that \( γ_{i,t_1} \in \{γ_{i,-1}, \ldots, γ_{t_1-1}\} \), where \( γ_{i,-1} = γ_{i,0} \) for odd \( i \) and \( γ_{i,-1} = γ_{i,k} \) for even \( i \). Let \(-1 \leq t_0 \leq t_1 \) be such that \( γ_{i,t_1} = γ_{i,t_0} \). By the assumption of \( γ_i \) being tangle-free, \( C = (γ_{i,t_0}, \ldots, γ_{i,t_1}) \) is the only directed cycle visited by \( γ_i \). The last important time, say \( (i, t_1) \), before \( (i, t_1) \) is called the short cycling time. We denote by \( t_2 \) the next time after \( (i, t_1) \) that is not an edge of \( C \) (\( γ_i \) circles around \( C \) between times \( (i, t_0) \) and \( (i, t_2) \)). We modify the mark of the short cycling time as \((γ_{i,t_1}, γ_{t_1}, γ_{t_1-1}, t_3)\) where \( τ = (i, t_3) \in T_i \cup \{(i, k)\} \) is the next time after \( (i, t_2) \) that \( e_τ \) will not be a tree edge (by convention \( τ = (i, k) \) if \( γ_i \) remains on the tree). Note that that this \((i, t_3)\) is the last important time, all steps to close the cycle are on tree edges. It follows that the pair \((γ_{i,t_3}, γ_{t_3-1})\) determines \( γ_{i,t_0} \). Important times \( (i, t) \) with \( 0 \leq t < t_1 \) or \( t_2 \leq t < k - 1 \) are called long cycling times. The other important times are superfluous. The key observation is that for each \( 1 \leq i \leq 2m \), the number of long cycling times \((i, t)\) is bounded by \( g - 1 \) (since there is at most one cycle, no edge of \( γ_i \) can be seen twice outside those of \( C \), \(-1 \) coming from the fact that the short cycling time is an important time). Now consider the case where the \( i \)-th path does...
not contain a cycle, then all important times are called long cycling times and their number is bounded by $g$.

We can reconstruct a canonical path $\gamma \in W_{k,m}$, from the sole positions of the short and long cycling times and their marks. This our second encoding. For each $i$, there are at most $(k+1)^2g$ ways to position the short and long cycling times of $T_i$, $s^2$ possibilities for the mark of a long cycling time and $s^2k$ possibilities for the mark of a short cycling time. We deduce that

$$|W_{k,m}(s,a)| \leq (k+1)^2g(s^2)^2m(g-1)(s^2k)^2m.$$  

Since $s \leq 2km$, the conclusion follows. \hfill \Box

Our second lemma bounds the contributions of paths in each equivalence class. This lemma is the new main technical difference of this section with [14].

**Lemma 15.3** Let $\gamma \in W_{k,m}$ such that $|V_\gamma| = s$ and $|E_\gamma| = a$. We have

$$\sum_{\gamma':\gamma'\sim\gamma} \prod_{e \in E_{\gamma'}} |P_e|^{m_e} \leq n^{-2km+s} K^{km-a} K^{6(a-s)+8m} \rho^{km}.$$  

**Proof** We first express the product of entries of $P$ in terms of the matrix $Q$:

$$\prod_{e \in E_\gamma} |P_e|^{m_e} = \prod_{e \in E_\gamma} |P_e|^{m_e-2} P_e^2 \leq \left( \frac{\sqrt{\rho K}}{n} \right)^{2km-2a} \prod_{e \in E_\gamma} \frac{Q_e}{n},$$  

where we have used 15.3, $m_e \geq 2$ and $|P_{xy}| = \sqrt{\rho K}/n$.

The statement of the lemma immediately follows from the claim:

$$\sum_{\gamma':\gamma'\sim\gamma} \prod_{e \in E_{\gamma'}} Q_e \leq \rho^a n^{r-a} K^{6(a-s)+8m}$$  

Indeed, let us check 15.5. Let us define the degree of a vertex $x$ in $V_\gamma$ as the sum of in-degrees and out-degrees: $\sum_{i,t} (1_{i,t-1=x} + 1_{i,t=x})$. Let $s_k$ and $s \geq k$ be the set of vertices of degree $k$ and at least $k$. We have

$$s_1 + s_2 + s \geq 3 = s \quad \text{and} \quad s_1 + 2s_2 + 3s \geq 3 \leq \sum_k ks_k = 2a.$$  

Subtracting the right-hand side to twice the left-hand side, we find

$$s \geq 3 \leq 2(a-s) + s_1 \leq 2(a-s) + 2m.$$  

The bound $s_1 \leq 2m$ comes from the fact that only the vertices $\gamma_{i,0}$ and $\gamma_{i,k}$, with $i \in [2m]$, can be of degree 1. Indeed, other vertices are of degree at least 2: for $1 \leq t \leq k$, $\gamma_{i,t}$ has in-degree at least 1 and out-degree at least 1.
Consider the subset of vertices \( \hat{V}_\gamma \subset V_\gamma \) which are of degree at least 3 or are among the extremes vertices \( \gamma_i, 0, \gamma_i, k, i \in [2m] \). In particular, \( V_\gamma \setminus \hat{V}_\gamma \) contains only vertices of degree 2. From what precedes
\[
\hat{s} = |\hat{V}_\gamma| \leq 2(a - s) + 4m. \tag{15.6}
\]
We may partition the edges of \( E_\gamma \) into \( \hat{a} \) sequences of edges of the form, for \( 1 \leq j \leq \hat{a} \),
\[
\hat{e}_j = (e_{j,1}, \ldots, e_{j,q_j}), \text{ with } e_{j,t} = (x_{j,t-1}, x_{j,t}) \in E_\gamma, x_{j,0}, x_{j,q_j} \text{ in } \hat{V}_\gamma \text{ and } x_{j,t} \not\in \hat{V}_\gamma \text{ for } 1 \leq t \leq q_j - 1.
\]
By construction
\[
\sum_{j=1}^{\hat{a}} q_j = a. \tag{15.7}
\]
We consider the directed graph \( \hat{G}_\gamma \) on the vertex set \( \hat{V}_\gamma \) whose \( \hat{a} \) edges are, for \( 1 \leq j \leq \hat{a} \),
\[
(x_{j,0}, x_{j,q_j}) \text{ (this is a multi-graph: if two sequences } \hat{e}_j \text{ and } \hat{e}_i, i \neq j, \text{ have the same extreme vertices, it creates two edges). It is straightforward to check that this operation preserves the genus:}
\[
a - s = \hat{a} - \hat{s}. \tag{15.8}
\]
For ease of notation, let \( y_1, \cdots, y_{\hat{s}} \) be the elements of \( \hat{V}_\gamma \). Let \( a_j \) and \( b_j \) the indices such that \( x_{j,0} = y_{a_j} \) and \( x_{j,q_j} = y_{b_j} \). Summing over all possible vertices, we get
\[
\sum_{\gamma': \gamma' \sim \gamma} \prod_{e \in E_{\gamma'}} Q_e \leq \sum_{(y_1, \cdots, y_{\hat{s}}) \in [n]^{\hat{s}}} \prod_{j=1}^{\hat{a}} Q_{y_{a_j} y_{b_j}}^{q_j},
\]
where we have used that
\[
\sum_{(x_{j,1}, \cdots, x_{j,q_j-1})} \prod_{t=1}^{q_j} Q_{x_{j,t-1} x_{j,t}}^{q_j} = Q_{x_{j,0}, x_{j,q_j}}^{q_j}.
\]
We apply (10.4) and find
\[
\sum_{\gamma': \gamma' \sim \gamma} \prod_{e \in E_{\gamma'}} Q_e \leq \sum_{(y_1, \cdots, y_{\hat{s}}) \in [n]^{\hat{s}}} \prod_{j=1}^{\hat{a}} \left( \frac{K^2 \rho q_j}{n} \right).
\]
Using 15.6–15.8, we have \( \hat{a} \leq 3(a - s) + 4m \) and, from 15.7, Eq. (15.5) follows. \( \square \)

We are ready for the proof of Lemma 15.1.
Proof of Lemma 15.1} Note that $W_{k,m}(s,a)$ is empty unless $0 \leq s - 1 \leq a \leq km$ (since each edge has multiplicity at least 2, we have $2|E_\gamma| \leq 2km$ from 15.3). From 15.4, we get

$$E\|A^{(k)}\|^{2m} \leq \sum_{a=1}^{km} \sum_{s=1}^{a+1} \left(\frac{n}{d}\right)^{2km-a} \max_{\gamma \in W_{k,m}(s,a)} \sum |P_e|^{me}.$$ 

Using Lemmas 15.2 and 15.3, we arrive at

$$E\|A^{(k)}\|^{2m} \leq n \sum_{a=1}^{km} \sum_{g=0}^{\infty} d^{-2km} (2km)^{6mg+2m} n^{-g} K^{km-a} K^{-6g+8m-6} \rho^{km},$$

where we have performed the change of $s \rightarrow g = a - s + 1$ and used $\vartheta_2 = \sqrt{\rho/d}$.

Recall $k \leq \ln(n)$. We take $m = \lceil \ln(n/K^6)/(12 \ln(\ln(n))) \rceil$. If $n \geq c K^6$ for some universal constant $c$, we find that

$$\frac{K^6(2km)^{6m}}{n} \leq \frac{1}{2}.$$ 

We deduce that

$$E\|A^{(k)}\|^{2m} \leq n \vartheta_2^{2km} (1 + \frac{K}{d})^{km} (2km)^{2m} K^{8m} (2km).$$

For our choice of $m$, $2km \leq (\ln(n))^2 / \ln(\ln(n))$ and, if $n \geq K^{42}$, then $n^{1/(2m)} \leq \ln(n)^7$. Since $\sqrt{K/d} = \vartheta_1$, the conclusion follows easily.

\begin{align*}
15.2 \textbf{Norm of } R^{(\ell)}_k \\
\end{align*}

In this section, we prove 13.11.

Lemma 15.4 There exists a universal constant $c \geq 3$ such that for all integers $1 \leq k \leq \ell \leq \ln(n)$ and $n \geq c$,

$$\left( E \left\| R^{(\ell)}_k \right\|^{2m} \right)^{\frac{1}{2m}} \leq \frac{d}{n} \ln(n)^{23} L\ell,$$

where $m = \ln(n)/(24 \ln(\ln(n)))$.

The proof follows from the same line than the proof of Lemma 15.1. It is also essentially contained in [14]. To avoid repetitions, we only focus on the main differences with the proof of Lemma 15.1. The computation leading to 15.1 gives
\[
E \| R_k^{(\ell)} \|^{2m} \leq \left( \frac{n}{d} \right)^{2(\ell - 1)m} \sum_{\gamma \in W'_{\ell,m}} E \prod_{i=1}^{2m} \prod_{t=1}^{\ell} P_{\gamma_{i,t-1} \gamma_{i,t}} \left( 1_{t<k} M_{\gamma_{i,t-1} \gamma_{i,t}} + 1_{t=k} + 1_{t>k} M_{\gamma_{i,t-1} \gamma_{i,t}} \right),
\]

where \( W'_{\ell,m} \) is the set of \( \gamma = (\gamma_1, \ldots, \gamma_{2m}) \) with \( \gamma_i = (\gamma_{i,0}, \ldots, \gamma_{i,\ell}) \notin F^\ell \), \( (\gamma_{i,0}, \ldots, \gamma_{i,k-1}) \in F^{k-1}, (\gamma_{i,k+1}, \ldots, \gamma_{i,\ell}) \in F^{\ell-k} \) and the boundary conditions: for all \( i \in [m] \),

\[
\gamma_{2i,0} = \gamma_{2i+1,0} \quad \text{and} \quad \gamma_{2i-1,\ell} = \gamma_{2i,\ell}
\]

with \( \gamma_{2m+1} = \gamma_1 \).

We associate to an element \( \gamma \in W'_{\ell,m} \) the directed graph \( G'_\gamma = (V'_\gamma, E'_\gamma) \) with vertices \( V'_\gamma = \{ \gamma_{i,t} : 1 \leq i \leq 2m, 0 \leq t \leq \ell \} \) and edge set \( E'_\gamma = \{ (\gamma_{i,t-1}, \gamma_{i,t}) : 1 \leq i \leq 2m, 1 \leq t \leq \ell, t \neq k \} \). The graph \( G'_\gamma \) is not necessarily weakly connected (since \( E'_\gamma \) does not contain the edges \( (\gamma_{i,k-1}, \gamma_{i,k}) \)). However, we have the following observation.

**Lemma 15.5** If \( \gamma \) is as above then each connected component of \( G'_\gamma \) contains a cycle. In particular, \( |E'_\gamma| \geq |V'_\gamma| \).

**Proof** By recursion, it is then enough to check that each connected component of \( G'_{\gamma_{i_1}} \) contains a cycle. By assumption, \( \gamma_i = (\gamma_{i,0}, \ldots, \gamma_{i,\ell}) \notin F^\ell \) contains two distinct cycles. Up to recomposing a new cycle, we may assume without loss of generality that the edge \( (\gamma_{i,k-1}, \gamma_{i,k}) \) is in zero or one of the two cycles. If it is in one of them, then the graph \( G'_{\gamma_{i_1}} \) is weakly connected and it contains the other cycle. Assume now that \( (\gamma_{i,k-1}, \gamma_{i,k}) \) is in none of the two cycles. If \( G'_{\gamma_{i_1}} \) is weakly connected, there is nothing to prove. If \( G'_{\gamma_{i_1}} \) is not weakly connected, then the two connected components are the vertices of \((\gamma_{i,0}, \ldots, \gamma_{i,k-1})\) and \((\gamma_{i,k}, \ldots, \gamma_{i,\ell})\). Since these two paths are tangle-free, each must contain exactly one of the two cycles and the statement follows. \( \square \)

Using the independence of the entries of \( M \) and \( |P_{xy}| \leq L/n \),

\[
E \prod_{i=1}^{2m} \prod_{t=1}^{\ell} P_{\gamma_{i,t-1} \gamma_{i,t}} \left( 1_{t<k} M_{\gamma_{i,t-1} \gamma_{i,t}} + 1_{t=k} + 1_{t>k} M_{\gamma_{i,t-1} \gamma_{i,t}} \right) \leq \left( \frac{d}{n} \right)^{|E'_\gamma|} \left( \frac{L}{n} \right)^{2\ell m}.
\]

We thus obtain that

\[
E \| R_k^{(\ell)} \|^{2m} \leq \left( \frac{d}{n} \right)^{2m} \left( \frac{L}{d} \right)^{2\ell m} \sum_{\gamma \in W'_{\ell,m}} \left( \frac{d}{n} \right)^{|E'_\gamma|},
\]

We introduce the equivalence class in \( W'_{\ell,m} \), we write \( \gamma \sim \gamma' \) if there exists a permutation \( \sigma \in S_n \) such that \( \gamma' = \sigma \circ \gamma \), where \( \sigma \) acts on \( \gamma \) by mapping \( \gamma_{i,t} \) to \( \sigma(\gamma_{i,t}) \). We denote by \( \mathcal{W}'_{\ell,m}(s,a) \) the set of equivalence classes such that \( |V'_\gamma| = s \)
and $|E'_γ| = a$. From Lemma 15.5, $W_{k,m}(s, a)$ is empty unless $a \geq s$. We have the following estimate on $W'_{ℓ,m}(s, a)$.

**Lemma 15.6** Let $a \geq s \geq 1$ be integers. We have

$$|W'_{ℓ,m}(s, a)| \leq (2ℓm)^{12m(a−s)+22m}.$$  

**Proof** A proof is contained in [14, Lemma 18]. We give a proof for the reader convenience. We order the sets $T' = \{(i, t) : 1 \leq i \leq 2m, 1 \leq t \leq ℓ, t \neq k\}$ and $T = \{(i, t) : 1 \leq i \leq 2m, 1 \leq t \leq ℓ − 1\}$ with the lexicographic order. We think of $T$ and $T'$ as times. If $τ ∈ T'$, we denote $τ^−$ the largest element in $T'$ smaller than $τ$. By convention, $(1, 0)^− = (1, −1)$. For $τ ∈ T$, we define $e_τ$ and $γ_τ$ as in Lemma 15.2 and we say $γ ∈ W'_{ℓ,m}$ is canonical if $V''_γ = \{1, \ldots, |V''_γ|\}$, $γ_{1,1} = 1$ and vertices are first visited in order. We aim for an upper bound on the number of canonical paths in $W'_{ℓ,m}$ with $|V''_γ| = s$ and $|E'_γ| = a$ by designing an injective map.

We define a sequence of growing sub-forests $(F_τ)_{τ ∈ T'}$ of $G'_γ$ as follows. We start with $F_{1,1}$, the trivial graph with no edge and a vertex $γ_{1,1} = 1$ as unique vertex. For $τ ∈ T'$, we say that $τ$ is a first time, if adding $e_τ$ to $F_τ$ does not create a weak cycle. If $τ$ is a first time the edge $e_τ$ is called a tree edge and we define $F_τ$ as the union of $e_τ$ and $F_τ$. Otherwise, $F_τ = F_τ$. We set $F = F_{2m,1}$. By construction, the set of tree edges is a sub-graph of $G'_γ$, with no weak cycle and vertex set $V''_γ$. Moreover, the weak connected components of $G'_γ$ and $F$ are equal. We call the other edges of $G'_γ$ the excess edges. In each weak connected component of $G'_γ$ there are at most $g = a − s + 1$ excess edges. Indeed, if $a', s'$ are the numbers of directed edges and vertices of a connected component, then there are $a' − s' + 1$ excess edges in this connected component. However by Lemma 15.5, $a − s' \leq a − s$. If $e_τ$ is an excess edge, we say that $τ ∈ T'$ is an important time. Other times in $T'$ are tree times (visit of a tree edge which has been seen before).

Let $k_i = k$ for odd $i$ and $k_i = k − ℓ + 1$ for even $i$. We define the sets $T^1_i = \{(i, t) : 0 \leq t \leq k_i − 1\}$ and $T^2_i = \{(i, t) : k_i \leq t \leq ℓ\}$. For each $i$, there could be a special first time $(i, t) ∈ T^2_i$, called the merging time, such that a connected component of $F_τ$ merges into a connected component of $F_i,k_i − 1$ by the addition of $e_{i,t}$.

The sets $T^ε_i$ are composed by the successive repetitions of $(i)$ a sequence of the tree times (possibly empty), $(ii)$ a sequence of first times (possibly empty), $(iii)$ an important time or the merging time. We mark the important and merging times (for $ε = 2$) $(i, t) ∈ T^ε_i$ by the vector $(y_{i,t}, y_{τ−1})$, where $τ ∈ T^ε_i \cup \{(i, k_i)\}$ is the next time that $e_τ$ will not be a tree edge (by convention $τ = (i, k_i)$ if $T^ε_i$ only contains tree times after $(i, t)$). We can reconstruct a canonical path $γ ∈ W'_{ℓ,m}$, from the positions of the merging and important times and their marks.

We refine this encoding by partitioning important times into three categories, short cycling, long cycling and superfluous times exactly as done in Lemma 15.2, except that there are short and long cycling times for each $i$ and $ε ∈ \{1, 2\}$ in the sequence $T^ε_i$. There are either 0 short cycling times and at most $g$ long cycling times, or 1 short cycling time and at most $g − 1$ long cycling time (because in each connected component of $G'_γ$, there are at most $g$ excess edges).
We can reconstruct a canonical path $\gamma \in W'_{\ell,m}$, from the positions of the merging, short and long cycling times and their marks. There are at most $\ell^2 m$ ways to position the merging times. For each $i, \varepsilon$, there are at most $\ell g$ ways to position the short and long cycling times of $T^\varepsilon_i$, $s^2$ possibilities for the marks of a merging or long cycling time and $s^2 k$ possibilities for the marks of a short cycling time. We deduce that

$$|W'_{\ell,m}(s, a)| \leq \ell^{4m+2m} (s^2)^{4m(g-1)+2m} (s^2 k)^{4m}. $$

Since $s \leq 2 \ell m$, the conclusion follows. $\square$

We are ready for the proof of Lemma 15.4.

**Proof of Lemma 15.4** There are $n(n-1) \cdots (n-s+1)$ elements of $W'_{\ell,m}$ in an equivalence class in $W'_{\ell,m}(s, a)$. From 15.10 and Lemma 15.6 we get

$$E\|R_k^{(\ell)}\|^{2m} \leq \left( \frac{d}{n} \right)^{2m} \left( \frac{L}{d} \right)^{2\ell m} \sum_{a=1}^{\ell m} \sum_{s=1}^{a} n^s (2\ell m)^{12m(a-s)+22m} \left( \frac{d}{n} \right)^{a}. $$

We perform the change of variable $s \rightarrow p = a - s$:

$$E\|R_k^{(\ell)}\|^{2m} \leq \left( \frac{d}{n} \right)^{2m} \left( \frac{L}{d} \right)^{2\ell m} (2\ell m)^{22m} \sum_{a=1}^{\ell m} d^a \sum_{p=0}^{\infty} \left( \frac{(2\ell m)^{12m}}{n} \right)^p. $$

Recall $\ell \leq \ln(n)$ and $d \geq 1$. We take $m = \lceil \ln(n)/(24 \ln(\ln(n))) \rceil$. If $n \geq c_1$ for some universal constant $c_1$, we find that

$$\frac{(2km)^{12m}}{n} \leq \frac{1}{2}. $$

We deduce that

$$E\|R_k^{(\ell)}\|^{2m} \leq \left( \frac{d}{n} \right)^{2m} L^{2\ell m} (2\ell m)^{22m} (2\ell m). $$

For our choice of $m$, $2\ell m \leq \ln(n)^2$. The conclusion follows easily. $\square$

**16 Proof of Theorem 5.4 and Corollary 5.5**

The proofs of Theorem 5.4 follows the same line than the proof of Theorem 5.4. In this section we explain the differences.
16.1 Main Technical Result

We will assume without loss of generality that \( \mu_1 = 1 \). For ease of notation, we set in this section \( d = \bar{d}, \vartheta = \bar{\vartheta}, M = \bar{M} \) and \( D = \bar{d} \lor 1.01 \). Note that since \( d \geq 1 \), we have \( D \leq 2d \).

We will assume in the sequel without loss of generality that

\[
\log(n) \geq 8 \log(d). \tag{16.1}
\]

Also, from 10.3 \( \rho \geq \mu_1^2 = 1 \). In particular \( \vartheta \geq \frac{1}{\sqrt{d}} \) and \( \tau_0 - 1 \leq \sqrt{d} \). If \( \ell_0 \) follows that in the statement of Theorem 5.4, we have

\[
\tau_0 - 2\ell \leq d^{\ell} \leq n^{1/8}. \tag{16.2}
\]

We set \( E_n = [n]^2 = \{(x, y) : x, y \in [n]\} \). With a slight abuse of notation, we can identify the non-backtracking matrix with its natural extension on \( E_n \) defined as follows: for all \( e = (x, y) \in E_n, f = (a, b) \in E_n \),

\[
B_{e, f} = \frac{n}{d} M_e M_f 1(y = a) 1(b \neq x) P_f.
\]

The advantage of this new equivalent definition of \( B \) is that it is defined on the deterministic set \( E_n \) rather than on the random set \( E \). Similarly, if \( \varphi \) is a vector in \( \mathbb{R} \), we set for all \( e = (x, y) \in E_n \),

\[
\varphi^+(e) = \frac{1_{e \in E}}{\sqrt{d}} \varphi(y) \quad \text{and} \quad \varphi^-(e) = \frac{1_{e \in E}}{\sqrt{d}} \varphi(x) \tag{16.3}
\]

If \( r_0 \geq 1 \), we define \( \Phi^\pm = (\varphi_1^\pm, \ldots, \varphi_{r_0}^\pm) \) and \( \Sigma = \text{diag}(\mu_1, \ldots, \mu_{r_0}) \). The candidate left and right eigenvectors \( (u_1, \ldots, u_{r_0}) \) and \( (v_1, \ldots, v_{r_0}) \) are the columns of the matrices:

\[
U = B^\ell \Phi^+ \Sigma^{-\ell} \quad \text{and} \quad V = (B^*)^\ell \Delta \Phi^- \Sigma^{-\ell - 1},
\]

where \( \Delta \in \mathcal{M}_{E_n}(\mathbb{R}) \) is the diagonal matrix defined for all \( e \in E_n \) by

\[
\Delta_{e,e} = n P_e.
\]

We set \( S = U \Sigma^\ell V^* \) and introduce the vector spaces \( H = \text{vect}(v_1, \ldots, v_{r_0}) \) and \( H' = \text{vect}(u_1, \ldots, u_{r_0}) \).

Finally, if \( r_0 = 0 \), then \( S \) is simply set to be the zero matrix, and \( H, H' \) are the trivial vector spaces.

Then, we define the matrix \( \hat{\Gamma}^{(t)} = d(\Gamma^{(t+1)} - I_{r_0}) \in \mathcal{M}_{r}(\mathbb{R}) \): that is,

\[
\hat{\Gamma}^{(t)}_{i,j} = d \sum_{s=1}^{r+1} \frac{1, Q^s \varphi_i^s \ldots \varphi_j^s}{(\mu_i \mu_j d)^s}. \tag{16.4}
\]
The proof of Lemma 9.1 implies the following.

**Lemma 16.1** (Properties of $\hat{\Gamma}^{(t)}$) For any $t$, the matrix $\hat{\Gamma}^{(t)}$ is a semi-definite positive matrix with eigenvalues greater than $\sigma$, and with

$$\sigma \leq \|\hat{\Gamma}^{(t)}\| \leq d r_0 L^4 \frac{\tau_0^2 - \tau_0^{2(t+2)}}{1 - \tau_0^2}.$$

The following result is the analog of Theorem 9.2 for the non-backtracking operator $B$.

**Theorem 16.2** (Algebraic structure of $U$ and $V$) There are a universal constant $c > 0$ and an event with probability greater than $1 - cn^{-1/4}$ such that the following holds:

\[
\begin{align*}
\|V^* B^\ell U - \Sigma^\ell\| &\leq C_1 n^{-1/4} \tau_0^{2\ell} \delta^\ell & (16.5) \\
\|U^* V - I_0\| &\leq C_1 n^{-1/4} \tau_0^{2\ell} & (16.6) \\
\|\left(\Phi^+\right)^* U - I_0\| &\leq C_1 n^{-1/4} \tau_0^{2\ell} & (16.7) \\
\|\left(\Phi^+\right)^* V - I_0\| &\leq C_1 n^{-1/4} \tau_0^{2\ell} & (16.8) \\
\|U^* U - \Gamma^{(\ell)}\| &\leq C_1 n^{-1/4} \tau_0^{2\ell} & (16.9) \\
\|V^* V - \hat{\Gamma}^{(\ell)}\| &\leq C_1 n^{-1/4} \tau_0^{2\ell} & (16.10) \\
\|B^\ell P_{H^\perp}\| &\leq C_2 \delta^\ell & (16.11) \\
\|P_{H^\perp} B^\ell\| &\leq C_2 \delta^\ell, & (16.12)
\end{align*}
\]

where $C_1 = cr_0 b^2 L^4 \ln(n)^{7/2}$ and $C_2 = cd L^8 \bar{r} \ln(n)^{10}$.

The same statements holds, with the same constants, if we replace $\ell$ by $\ell' = \ell + 1$. Repeating the proof of Theorem 4.3 in Sect. 9.4 and Corollary 4.6, we obtain Theorem 5.4.

We now check that Corollary 5.5 holds. We start by a comment. We first observe the matrix identity

$$B^* \Delta J = \Delta J B$$

where $J$ is the matrix such that for all $(x, y) \in E$, we have $J \delta_{(x, y)} = \delta_{(y, x)}$. In particular, we find that $\Delta J \psi_i$ and $\psi'_i$ are proportional vectors: there exists $c > 0$ such that for all $(x, y) \in E$, $c \psi'_i(x, y) = n P_{x, y} \psi_i(y, x)$. If $\deg(i) = 1$, we have that $\psi_i(y) = 0$. If $\deg(y) \geq 2$, using the eigenvalue equation, we deduce that

$$\lambda_i \hat{\psi}_i(y) = \frac{1}{d(\deg(y) - 1)} \sum_{x: (x, y) \in E} \lambda_i \psi_i((x, y)) = \frac{1}{d} \sum_{x: (x, y) \in E} n P_{y, x} \psi_i((y, x)) = c \hat{\psi}_i(y).$$
The estimators \( \hat{\psi}_i \) and \( \hat{\psi}_i \) are thus very close (note however that if \( \text{deg}(i) = 1 \), we have that \( \hat{\psi}_i(y) = 0 \) but \( \hat{\psi}_i'(y) \) could be different from 0).

Let us now check that the conclusion of Corollary 5.5 holds. Arguing as in the proof of Corollary 4.6, it suffices to check the conclusion with our approximate eigenvectors \( u_i \) and \( v_i \) in place of the actual eigenvectors \( \psi_i \) and \( \psi_i' \). It then follows from a slight modification of Theorem 16.2 that, with probability at least 1 – \( cn^{-1/4} \), \(|\langle \hat{u}_i, \hat{u}_j \rangle - \delta_{i,j} \Gamma^{(e)}_{i,j} \| \leq C_1 n^{-1/4} T_0^{2\ell} \) and \(|\langle \hat{u}_i, \varphi_j \rangle - \delta_{i,j} \| \leq C_1 n^{-1/4} T_0^{2\ell} \). The same holds with the vectors \( \hat{v}_i \). The conclusion follows.

## 16.2 Near Eigenvectors: Claims 16.5–16.10

We denote by \( G = ([n], E) \) the undirected random graph whose adjacency matrix is \( M \): this is an Erdős-Rényi random graph with parameter \( d/n \) (with loops).

The first step is to extend the notion of graph neighborhood \( (G, x)_y \) to directed edge neighborhood. If \( e = (y, x) \in E_n \), we define \( (G, e) \), has the rooted graph \( G^e \) where \( G^e \) is the graph obtained from \( G \) by removing the edge \( \{x, y\} \) if it is in \( E \) otherwise \( G^e = G \). As explained in [14], Lemma 11.3 and Proposition 11.4 also hold for \( (G, e) \) in place of \( (G, x) \) and \( (T, x) = (T, e) \) is now the marked Poisson Galton-Watson tree with parameter \( d \) with root mark \( x \) if \( e = (y, x) \).

If \( (g, o) \in \mathcal{G}_s \), we write \( o \sim u \) if \( u \) and \( o \) share an edge and \( (g, (o, u)) \) is defined as above as the rooted graph rooted at \( u \) where the edge with \( o \) has been removed. The version of Theorem 12.5 that we will need is the following:

**Theorem 16.3** Let \( 0 \leq \kappa \leq 0.49 \) and \( h \) an integer as in (11.10). Let \( f : \mathcal{G}_s \to \mathbb{R} \) be a \((h + 1)\)-local function such that \( |f(g, o)| \leq \alpha \sum_{u \sim o} |(g, (o, u))_h|^{\beta} \) for some \( \alpha, \beta > 0 \). Then, for some universal constant \( c > 0 \), for any \( s \geq 1 \), with probability greater than 1 – \( n^{-s} \), we have

\[
\left| \sum_{x \in [n]} f(G, x) - E \sum_{x \in [n]} f(T, x) \right| \leq c e^\beta \alpha s^{5/2 + \beta} \ln(n)^{5/2 + \beta} Dn^{x(1 + \beta) + 1/2}.
\]

The proof is the same than the proof of Theorem 12.5. The only difference is at the first lines: we set \( \tilde{f}(g, o) = \alpha \sum_{u \sim o} |(g, (o, u))_h|^{\beta} / d \). Note that if an integer random variable \( N \) is independent of \( (Z_k) \) i.i.d. non-negative, then for any \( q \geq 1 \),

\[
E \left[ \left( \sum_{k=1}^{N} Z_k \right)^q \right] \leq E \left[ N^{q-1} \sum_{k=1}^{N} Z_k^q \right] = E[N^q]E[Z_1^q].
\]

Hence, by 11.8 (applied to \( t = h \) and \( t = 1 \)), we have for all \( p \geq 1 \) and \( x \in [n] \), for some universal constant \( c > 0 \),

\[
E[\tilde{f}(G, x)^{2p}] \leq c^{2p} (cp)^{2p(\beta + 1)} D^{2p(\beta h + 1)}.
\]

The rest is identical.
Theorems 16.3 and 12.2 are used to prove claims 16.5–16.10. As an illustration, we first check that (16.9) holds. The entry $(i, j)$ of $U^*U$ is equal to

$$(\mu_i \mu_j)^{-\ell} (B^\ell \varphi_i^+, B^\ell \varphi_j^+) = \sum_{x, y \in [n]} (\mu_i \mu_j)^{-\ell} \mathbf{1}_{[x, y] \in E} (B^\ell \varphi_i^+(y, x))(B^\ell \varphi_j^+(y, x)).$$

On the event that $G$ is $\ell$-tangle free, we have

$$(B^\ell \varphi_i^+(y, x))(B^\ell \varphi_j^+(y, x)) = f(G, (y, x))/d,$$

where

$$f(g, o) = 1_{(g, o)\ell} \text{ is tangle free } f_{1, \varphi_i, \ell}(g, o) f_{1, \varphi_j, \ell}(g, o).$$

We thus have on the event that $G$ is $\ell$-tangle free,

$$(\mu_i \mu_j)^{-\ell} (B^\ell \varphi_i^+, B^\ell \varphi_j^+) = \sum_{x \in [n]} \tilde{f}(G, x),$$

with

$$\tilde{f}(g, o) = (\mu_i \mu_j)^{-\ell} \frac{1}{d} \sum_{u \sim o} f(g, (o, u)),$$

and the sum is over all neighbors of $o$ (end vertex of an edge attached to $o$). The function $\tilde{f}(g, o)$ is $(\ell + 1)$-local and, from Lemma 12.1, it is bounded by

$$\frac{2b^2}{nd} |\mu_i \mu_j|^{-\ell} \mu^2 \sum_{u \sim o} |(g, (o, u))\ell|^2.$$ 

Then claim 16.9 follows by applying Theorem 16.3 and Theorem 12.2 - Eq. (12.7).

We now check that (16.10) holds. The entry $(i, j)$ of $V^*V$ is equal to $(\mu_i \mu_j)^{-\ell-1}$ times

$$\langle (B^*)^\ell \Delta \varphi_i^-, (B^*)^\ell \Delta \varphi_j^- \rangle = \sum_{x, y \in [n]} \mathbf{1}_{[x, y] \in E} ((B^*)^\ell \Delta \varphi_i^-)(x, y)((B^*)^\ell \Delta \varphi_j^-)(x, y).$$

Using the symmetry of $E$ and $P$, we find that on the event that $G$ is $\ell$-tangle free,

$$((B^*)^\ell \Delta \varphi_i^-)(x, y)((B^*)^\ell \Delta \varphi_j^-)(x, y) = n^2 P_{x, y} f(G, (y, x))/d = (n/d) Q_{x, y} f(G, (y, x)),$$

where as above

$$f(g, o) = 1_{(g, o)\ell} \text{ is tangle free } f_{1, \varphi_i, \ell}(g, o) f_{1, \varphi_j, \ell}(g, o).$$
We thus have on the event that $G$ is $\ell$-tangle free,

$$(\mu_i \mu_j)^{-\ell-1} \langle (B^\ell)^2 \Delta \varphi^-_i, (B^\ell)^2 \Delta \varphi^-_j \rangle = \sum_{x \in [n]} \tilde{f}(G, x),$$

with

$$\tilde{f}(g, o) = (\mu_i \mu_j)^{-\ell-1} \frac{n}{d} \sum_{u \sim o} Q_{(o), (u)} f(g, (o, u)),$$

and $\iota$ is the mark function. The function $\tilde{f}$ is $(\ell + 1)$-local and it is bounded by

$$\frac{2b^2L^2}{nd} |\mu_i \mu_j|^{-\ell} \sum_{u \sim o} |(g, (o, u))|_\ell^2,$$

where we have used that $\mu_1 = 1$ and $K = \max_{x, y} n Q_{x, y} \leq L^2 / \mu_1^2$ as in 10.2. Moreover, by Theorem 12.2 - Eq. (12.7), we have

$$\mathbb{E} \tilde{f}(T, x) = \sum_y \frac{Q_{x, y}}{\mu_i \mu_j} \sum_{s=0}^{\ell} Q^s \varphi^{i, j}(y) (\mu_i \mu_j d)^s = d \sum_{s=1}^{\ell+1} Q^s \varphi^{i, j}(x) (\mu_i \mu_j d)^s.$$

It follows that claim 16.9 follows by applying Theorem 12.5.

We now check that 16.6 holds. The entry $(i, j)$ of $U^* V$ is equal to $\mu_i^{-\ell} \mu_j^{-\ell-1}$ times

$$\langle B^\ell \varphi^+_i, (B^\ell)^2 \Delta \varphi^-_j \rangle = \sum_{x, y \in [n]} 1_{(x, y) \in E}(B^2 \varphi^+_i)(y, x)(\Delta \varphi^-_j)(y, x).$$

On the event that $G$ is $\ell$-tangle free, we have

$$(B^2 \varphi^+_i)(y, x)(\Delta \varphi^-_j)(y, x) = (n/d) P_{x, y} \varphi_j(y) f(G, (y, x)),$$

where

$$f(g, o) = 1_{(g, o)_{\ell}} \text{ is tangle free } f_{1, \psi_i, 2\ell}(g, o).$$

We thus have on the event that $G$ is $\ell$-tangle free,

$$\mu_i^{-\ell} \mu_j^{-\ell-1} \langle B^\ell \varphi^+_i, (B^\ell)^2 \Delta \varphi^-_j \rangle = \sum_{x \in [n]} \tilde{f}(G, x),$$

with

$$\tilde{f}(g, o) = \mu_i^{-\ell} \mu_j^{-\ell-1} \frac{n}{d} \sum_{u \sim o} P_{(u), (o)} \varphi_j(\iota(o)) f(g, (o, u)).$$
and $i$ is the mark function. The function $\tilde{f}(g, o)$ is $(\ell + 1)$-local and, from Lemma 12.1, it is bounded by $(2b^4/(nd)|\mu_i\mu_j|^{\ell-\ell}\delta^{2\ell}\sum_{u \sim o} |(g, (o, u))|_\ell$. Moreover by Theorem 12.2 - Eq. (12.6), we have

$$E\tilde{f}(T, x) = \mu_j \mu_j^{-\ell-1} \sum_{y} P_{y,x} \varphi_j(y) \varphi_i(x) = \mu_j \mu_j^{-\ell} \varphi_j(x) \varphi_i(x).$$

where at the second step, we use that $\varphi_j$ is a left eigenvector of $P$. Then claim 16.9 follows by applying Theorem 16.3 and using the orthogonality relation $\langle \varphi_i, \varphi_j \rangle = \delta_{i,j}$.

We leave the remaining inequalities 16.5, 16.7 and 16.8 to the reader.

### 16.3 Norm of the Restricted Matrix: Claims 16.11–16.12

We may perform essentially the same tangle-free decomposition as performed for the matrix $A$. We follow [14] where a tangle-free decomposition for non-backtracking matrices was previously performed. The following argument introduces a new simplification.

We define a *non-backtracking path* of length $k$ as a sequence $(x_0, \ldots, x_k)$ in $[n]$ such that $x_t \neq x_{t+1}$ for all $t \in [k - 1]$. For example, if $x \neq y$, $(x, x, y, y)$ is a non-backtracking path of length 3. For $e, f \in E_n = [n]^2$, the $F_{\ell+1}^e,f$ is the set of all non-backtracking paths $(x_0, \ldots, x_{\ell+1})$ such that $(x_0, x_1) = e$, $(x_{\ell}, x_{\ell+1}) = f$ and the graph of the path is tangle-free (see definition in Sect. 11.1 on page 50). More generally, $F^t$ denotes the set of all tangle-free non-backtracking paths of length $t$, whatever their endpoints.

If the graph $G$ is $(\ell + 1)$-tangle free then $B_\ell = \tilde{B}(\ell)$ where, for $e = (x, y)$ and $f = (a, b)$ in $E_n = [n]^2$,

$$B_{e,f}^{(\ell)} = \left(\frac{n}{d}\right)^\ell \sum_{P_{e,f}^{\ell+1}} M_e \prod_{t=1}^{\ell} P_{x_t,x_{t+1}} M_{x_t,x_{t+1}}.$$

We also define the matrices $M$ and $\tilde{B}(\ell)$, $\tilde{B}_{e,f}^{(\ell)}$ by $B_{e,f}^{(0)} = \tilde{B}_{e,f}^{(0)} = I_{e,f} \in E$, and

$$M_{x,y} = M_{x,y} - \frac{d}{n},$$

$$B_{e,f}^{(\ell)} = \left(\frac{n}{d}\right)^\ell \sum_{P_{e,f}^{\ell+1}} M_e \left(\prod_{t=1}^{\ell-1} P_{x_t,x_{t+1}} M_{x_t,x_{t+1}}\right) P_f M_f,$$

$$\tilde{B}_{e,f}^{(\ell)} = \left(\frac{n}{d}\right)^\ell \sum_{P_{e,f}^{\ell+1}} M_e \left(\prod_{t=1}^{\ell-1} P_{x_t,x_{t+1}} M_{x_t,x_{t+1}}\right) P_f M_f.$$

(16.13)

Note that the difference between $B^{(\ell)}$ and $\tilde{B}^{(\ell)}$ lies at the weight attached to the last edge.
We use the convention that the product over an emptyset is 1. For each 
\((x_0, \ldots, x_{\ell+1})\), we write the telescopic sum (13.8) for the indices \(t \in [\ell - 1]\), we get

\[
B^{(\ell)}_{e, f} = B_{e, f}^{(\ell)} + \sum_{k=1}^{\ell-1} \left( \frac{n}{d} \right)^{\ell} \sum_{e \in \mathcal{E}} \sum_{f \in \mathcal{F}} M_e \prod_{t=1}^{k-1} P_{x_t x_{t+1}} M_{x_t x_{t+1}} \left( \frac{d}{n} P_{x_k x_{k+1}} \right) \prod_{t=k+1}^{\ell} P_{x_t x_{t+1}} M_{x_t x_{t+1}}.
\]

(16.14)

We introduce the matrix in \(\mathcal{M}_{E_n}(\mathbb{R})\): for \(e = (x, y)\), \(f = (a, b)\),

\[H_{e, f} = 1_{y = a, x \neq b}\]

The matrix \(H\) is the unweighted non-backtracking matrix on \([n]^2\).

The above identity 16.14 can then be rewritten as the following identity in \(\mathcal{M}_{E_n}(\mathbb{R})\):

\[
B^{(\ell)} = B^{(\ell)} + \frac{1}{nd} \sum_{k=1}^{\ell-1} \left( \frac{n}{d} \right)^{\ell} B_{e, f}^{(k-1)} \left( H \Delta \right)^{2} B^{(\ell-k-1)} - \sum_{k=1}^{\ell-1} R_k^{(\ell)}
\]

(16.15)

where

\[
(R_k^{(\ell)}{e, f}) = \left( \frac{n}{d} \right)^{\ell} \sum_{e \in \mathcal{E}} \sum_{f \in \mathcal{F}} M_e \prod_{t=1}^{k-1} P_{x_t x_{t+1}} M_{x_t x_{t+1}} \prod_{t=k+1}^{\ell} P_{x_t x_{t+1}} M_{x_t x_{t+1}}
\]

The only difference with 16.3 is that we do not require that \(e \in E\). Now, as in [14], we write that for \(e = (x, y)\), \(f = (a, b)\)

\[
\frac{1}{n} (H \Delta H)_{e, f} = 1_{a \neq x, y \neq b} P_{ya}
\]

\[
= P_{ya} - 1_{\{a = x\} \cup \{y = b\}} P_{ya}
\]

\[
= d \sum_{j=1}^{\mu} \mu_j \tilde{\varphi}_j^+(e) \tilde{\varphi}_j^-(f) - d S_{e, f},
\]

\(\tilde{\varphi}_j^-(f) = \frac{\varphi(x)}{\sqrt{d}}\) and \(\tilde{\varphi}_j^+(e) = \frac{\varphi(y)}{\sqrt{d}}\).

(16.16)
where we have used the spectral decomposition $P = \mu_1 \varphi_1^* + \cdots + \mu_n \varphi_n^*$, used notation (16.3) and defined $S \in \mathcal{M}_{E_n}(\mathbb{R})$ as

$$S_{e.f} = 1_{\{a=x\}} \cup \{y=b\} \frac{P_y a}{d}.$$ 

For any unit vector $w$, we thus have

$$\frac{1}{nd} B^{(k-1)} (H \Delta)^2 B^{(\ell-k-1)} w = \sum_{j=1}^{n} \mu_j \hat{B}_{j}^{(k-1)} \varphi_j^* |\langle \varphi_j^- , B^{k} w \rangle|$$

$$- \hat{B}^{(k-1)} S \Delta B^{(\ell-k-1)} w.$$ 

We observe that on the image of $B$, $\varphi^-$ and $\hat{\varphi}^-$ coincide. Hence, in the above expression, we may replace $\hat{\Delta} \varphi^-_j$ by $\Delta \varphi^-_j$. From the triangle inequality,

$$\left| \frac{1}{nd} B^{(k-1)} (H \Delta)^2 B^{(\ell-k-1)} w \right| \leq \sum_{j=1}^{n} |\mu_j| |\hat{B}_{j}^{(k-1)} \varphi_j^*| |\langle \Delta \varphi^-_j , B^{(\ell-k-1)} w \rangle|$$

$$+ \| \hat{B}^{(k-1)} S \Delta B^{(\ell-k-1)} \|.$$ 

(16.17)

From the direct analog of Proposition 13.3 for $|\langle (B^*)^t \Delta \varphi^-_i , w \rangle|$, with probability at least $1 - cn^{2k-1}$, the following holds for any $t \leq \ell$ and $i \in [r_0]$ and $w \in H^\perp$:

$$|\langle \Delta \varphi^-_i , B^t w \rangle| \leq c L^2 \ell \vartheta^t.$$ 

From the direct analog of Proposition 13.2 for $|\langle (B^*)^t \Delta \varphi^-_i , w \rangle|$, with probability at least $1 - cn^{2k-1}$, for all $i \in [n] \setminus [r_0]$ and $t \leq \ell$, we have

$$|\langle (B^*)^t \Delta \varphi^-_i \rangle|^2 \leq \mu_i^{2t} \hat{\varphi}_i^t + c L^3 (\ln n)^{9/2} n^{3k/2 - 1/2} \vartheta^{2t}.$$ 

However, from Eqs. (10.7–10.2), we have $\mu_i^{2t} \hat{\varphi}_i^t \leq L^4 (t + 1) \vartheta^{2t}$. As a consequence, for some universal constant $c > 0$, for all $i \in [n] \setminus [r_0]$ and $t \leq \ell$,

$$|\langle \Delta \varphi^-_i , B^t w \rangle| \leq |w| |\langle (B^*)^t \Delta \varphi^-_i \rangle| \leq c L^4 \sqrt{\ell} \vartheta^t.$$ 

On the union of the two events, in (16.17) we find, for all $w \in H^\perp$:

$$\sum_{j=1}^{n} |\mu_j| |\hat{B}_{j}^{(k-1)} \varphi_j^*| |\langle \Delta \varphi^-_j , B^{(\ell-k-1)} w \rangle| \leq c L^4 \ell \sum_{j=1}^{n} |\mu_j| \vartheta^{\ell-k-1} |\hat{B}_{j}^{(k-1)} \varphi_j^*|.$$
Hence, we find from 16.15–16.17 that, for some $c > 0$, with probability at least $1 - cn^{2k-1}$

$$\| B^{\ell} P_{H'} \| \leq \| B^{(\ell)} \| + cL^4 \ell \sum_{k=1}^{\ell-1} \sum_{j=1}^{n} |\mu_j|^{\ell-k+1} |B^{(k-1)} \varphi_j^+|$$

$$+ \sum_{k=1}^{\ell-1} \| \tilde{B}^{(k-1)} S \Delta B^{(\ell-k-1)} \| + \sum_{k=1}^{\ell-1} \| R_k^{(\ell)} \|. \quad (16.18)$$

The next proposition gathers the necessary norm bounds.

**Proposition 16.4** There exists a universal constant $c > 0$ such that if $n \geq cK^{42}$, with probability at least $1 - 1/\sqrt{n}$, the following holds for any $k \in [\ell]$ and $j \in [n]$:

$$\| B^{(k)} \| \leq c \sqrt{d} \ln(n)^{10} K^4 \vartheta^k, \quad (16.19)$$

$$\| R_k^{(\ell)} \| \leq \frac{cd^{3/2}}{n} \ln(n)^{23} L^\ell, \quad (16.20)$$

$$|\tilde{B}^{(k)} \varphi_j^+| \leq cb \ln(n)^5 K \vartheta^k, \quad (16.21)$$

$$\| \tilde{B}^{(k-1)} S \Delta B^{(\ell-k-1)} \| \leq c \ln(n)^{19} K^4 \vartheta^\ell \frac{d^\ell}{\sqrt{n}}. \quad (16.22)$$

Before checking this proposition in the next section, let us check that it implies Claim 16.11. It suffices to use Propositions 16.4 and 16.18. We recall that $L = d \vartheta_1 \leq d \vartheta, \theta \geq 1/\sqrt{d}$ and $d^\ell \leq n^{1/8}$ from 16.2 and we arrive at Claim 16.11.

The proof of Claim 16.12 follows from the same argument by considering the transpose of 16.15. We omit the proof.

### 16.4 Proof of Proposition 16.4

In this section, we prove Proposition 16.4. It relies on similar ideas than the proof of Proposition 13.5.

**Claim (16.19)**

From Markov inequality, the following lemma implies (16.19).

**Lemma 16.5** There exists a universal constant $c \geq 3$ such that for all integers $1 \leq k \leq \ln(n)$ and $n \geq cK^{42}$,

$$\left( E \left\{ \| B^{(k)} \|^{2m} \right\} \right)^{\frac{1}{2m}} \leq \ln(n)^9 \sqrt{d} K^4 \vartheta^k,$$

where $m = \ln(n/K^6)/(12 \ln(\ln(n)))$. 

\[ \textcircled{1} \text{ Springer} \]
The proof of this lemma is almost the same than the proof of Lemma 15.1. The computation leading to 15.1 gives
\[
\mathbb{E} \| B^{(k)} \|^{2m} \leq \left( \frac{n}{d} \right)^{2km} \sum_{\gamma} \mathbb{E} \prod_{i=1}^{2m} M_{\gamma_i,0,\gamma_{i+1}} \left( \prod_{t=1}^{k-1} P_{\gamma_t,\gamma_{t+1}} M_{\gamma_t,\gamma_{t+1}} \right) P_{\gamma_{2k},\gamma_{2k+1}} M_{\gamma_{2k},\gamma_{2k+1}} ,
\]
(16.23)
where the sum is over all \( \gamma = (\gamma_0, \ldots, \gamma_{2m}) \) with \( \gamma_i = (\gamma_{i,0}, \ldots, \gamma_{i,k+1}) \in F^{k+1} \) and the boundary conditions: for all \( i \in [m] \),
\[
(\gamma_{2i,0}, \gamma_{2i,1}) = (\gamma_{2i+1,0}, \gamma_{2i+1,1}) \quad \text{and} \quad (\gamma_{2i-1,k}, \gamma_{2i-1,k+1}) = (\gamma_{2i,k}, \gamma_{2i,k+1})
\]
with \( \gamma_{2m+1} = \gamma_1 \).
We associate to an element \( \gamma \) as above, an undirected graph \( G_\gamma = (V_\gamma, E_\gamma) \) with vertices \( V_\gamma = \{\gamma_i : 1 \leq i \leq 2m, 0 \leq t \leq k+1 \} \) and edge set \( E_\gamma = \{\{\gamma_i, \gamma_i+1 \} : 1 \leq i \leq 2m, 0 \leq t \leq k \} \). From the above boundary conditions, the graph \( G_\gamma \) is simply connected. In particular, \( |E_\gamma| - |V_\gamma| + 1 \geq 0 \). Each edge \( e \in E_\gamma \) has a multiplicity \( m_e \) defined as the number of times it is visited by \( \gamma \). By construction,
\[
\sum_{e \in E_\gamma} m_e = 2(k+1)m.
\]
(16.24)
We also define \( \tilde{m}_e \leq m_e \) as the number of visits of \( e \in E \) for some \( \{\gamma_i, t, \gamma_{i+1} \} \) with \( t \in [k] \) (we exclude \( t = 0 \)).
We may now estimate the expectation on the right-hand side of (16.23). Recall that the random variables \( M_{xy} = (M_{xy} - d/n) \), \( x, y \), are i.i.d., centered, bounded by 1 and with variance \( (d/n) - (d/n)^2 \). It follows that for any \( p \geq 1 \), \( \mathbb{E} |M_{xy}|^p \leq d/n \). Note also that for all \( p \geq 1 \), \( \mathbb{E} M_{xy}^p = d/n \). We deduce that
\[
\mathbb{E} \prod_{i=1}^{2m} M_{\gamma_i,0,\gamma_{i+1}} \left( \prod_{t=1}^{k-1} P_{\gamma_t,\gamma_{t+1}} M_{\gamma_t,\gamma_{t+1}} \right) M_{\gamma_{2k},\gamma_{2k+1}} \leq \left( \frac{d}{n} \right)^{|E_\gamma|} \prod_{e \in E_\gamma} |P_e|^{\tilde{m}_e}.
\]
Moreover, the above expectation is zero unless all edges have multiplicity at least 2 (note that \( M_{xy} \) is not centered but from the boundary condition, the edges \( \{\gamma_0, \gamma_1 \} \) and \( \{\gamma_k, \gamma_{k+1} \} \) have multiplicity at least 2). From (15.1), we thus obtain that
\[
\mathbb{E} \| B^{(k)} \|^{2m} \leq \sum_{\gamma \in \bar{W}_{k+1,m}} \left( \frac{n}{d} \right)^{2km - |E_\gamma|} \prod_{e \in E_\gamma} |P_e|^{\tilde{m}_e},
\]
(16.25)
where \( \bar{W}_{k+1,m} \) is the set of paths \( \gamma \) as above such that each edge of \( E_\gamma \) is visited at least twice.

The right-hand side of Eq. (16.25) is very close to the right-hand side of Eq. (15.4) for \( k' = k+1 \). It can be analyzed with the same method. We define \( \tilde{W}_{k,m}(s, a) \) as the
set of equivalence classes of paths in $\bar{W}_{k,m}$ with $|V_\gamma| = s$ and $|E_\gamma| = a$. We have that $\bar{W}_{k,m}(s, a)$ is empty unless $a - s + 1 \geq 0$. The conclusion of Lemma 15.2 also holds for $\bar{W}_{k,m}(s, a)$.

**Lemma 16.6** Let $a, s \geq 1$ be integers such that $a - s + 1 \geq 0$. We have

$$|\bar{W}_{k,m}(s, a)| \leq (2km)^{6m(a-s+1)+2m}.$$  

**Proof** We may repeat the proof Lemma 16.6. The new property that we use that there exists a unique non-backtracking path between two vertices in an (undirected) tree. \(\Box\)

Our second lemma bounds the contributions of paths in each equivalence class. Due to the presence of the multiplicities $\tilde{m}_e$ instead of standard multiplicities $m_e$, there is an extra factor in the analog of Lemma 15.3.

**Lemma 16.7** Let $\gamma \in \bar{W}_{k+1,m}$ such that $|V_\gamma| = s$ and $|E_\gamma| = a$. We have

$$\sum_{\gamma' : \gamma' \sim \gamma} \prod_{e \in E_{\gamma'}} |P_e|^{\tilde{m}_e} \leq n^{-2km+s} K^{km-a} K^{6(a-s)+8m} \rho^{km}. $$

**Proof** We let $\hat{\gamma} = (\hat{\gamma}_1, \ldots, \hat{\gamma}_{2m})$ be obtained from $\gamma$ by setting $\hat{\gamma}_i = (\tilde{\gamma}_{i,1}, \ldots, \tilde{\gamma}_{i,k+1})$ (we remove the first step for each $i$). Due to the boundary condition, the associated graph $G_{\hat{\gamma}}$ is connected and

$$\tilde{a} - \tilde{s} \leq a - s, \quad (16.26)$$

with $\tilde{a} = |E_{\hat{\gamma}}|$ and $\tilde{s} = |V_{\hat{\gamma}}|$. Note also $\tilde{m}_e$ is the number of visits of $e$ in $\hat{\gamma}$. We have that

$$\sum_{e \in E_{\hat{\gamma}}} \tilde{m}_e = \sum_{e \in E_{\hat{\gamma}}} \tilde{m}_e = 2km.$$ 

Let $H \subset E_{\hat{\gamma}}$ be the subset of edges $e \in E_{\hat{\gamma}}$ such that $\tilde{m}_e = 1$. We set $|H| = h$. We have

$$\prod_{e \in E_{\hat{\gamma}}} |P_e|^{\tilde{m}_e} = \prod_{e \in E_{\hat{\gamma}}} |P_e|^{\tilde{m}_e-2} P_e^2 \leq \left( \frac{\sqrt{\rho} K}{n} \right)^{2km-2\tilde{a}} \left( \frac{\sqrt{\rho} K}{n} \right)^{2h} \prod_{e \in E_{\hat{\gamma}} \setminus H} Q_e \frac{Q_e}{n},$$  

where we have used $\sum_{e \notin H} \tilde{m}_e = 2km - h$ and $|P_{xy}| = \sqrt{\rho K}/n$.

As in the proof of Lemma 15.3, we consider the graph $\hat{G}_{\hat{\gamma}}$ obtained from $G_{\hat{\gamma}}$ by gluing vertices of degree 2 while keeping the extreme vertices $\hat{\gamma}_{i,1}, \hat{\gamma}_{i,k+1}, \hat{\gamma} \in [2m]$. Note that due the boundary conditions, the edges in $H$ are not modified by this operation: they remain edges in $\hat{G}_{\hat{\gamma}}$. The proof of 15.5 applied to the graph $\hat{G}_{\hat{\gamma}}$ gives

$$\sum_{\gamma' : \gamma' \sim \gamma} \prod_{e \in E_{\hat{\gamma}} \setminus H} Q_e \leq \rho^{\tilde{a} - h} n^{-\tilde{a} + h} K^{6(\tilde{a} - \tilde{s})+8m}. $$

\(\Box\) Springer
Therefore, we get from 16.27 and 16.26 that
\[
\sum_{\gamma'} \prod_{e \in E_{\tilde{\gamma}}} |P_e|^{\tilde{m}_e} \leq n^{-2km+s} K^{km-a} K^{6(a-s)+8m} \rho^{km} K^h.
\]
Since \( K \geq 1 \), this concludes the proof. \( \square \)

We are ready for the proof of Lemma 16.5.

**Proof of Lemma 16.5** Note that \( \tilde{W}_{k+1,m}(s, a) \) is empty unless \( 0 \leq s - 1 \leq a \leq (k+1)m \) (since each edge has multiplicity at least 2). From (16.25), we get
\[
\mathbb{E} [\|B^{(k)}\|^2] \leq \left( \frac{n}{d} \right)^{2km-a} |\tilde{W}_{k+1,m}(s, a)| \max_{\gamma' \in \tilde{W}_{k+1,m}(s,a)} \sum_{\gamma' : \gamma' \sim \gamma} \prod_{e \in E_{\gamma'}} |P_e|^{\tilde{m}_e}.
\]
We use Lemmas 16.6 and 16.7 and follow the computation of Lemma 15.1. We find
\[
\mathbb{E} [\|B^{(k)}\|^2] \leq n(2km)^{2m} K^{8m-6} \phi_2^{2km} \sum_{a=1}^{(k+1)m} \left( \frac{K}{d} \right)^{km-a} \sum_{g=1}^{\infty} \left( \frac{K^6 (2(k+1)m)^{6m}}{n} \right)^g.
\]
The conclusion follows easily. The conclusion follows easily. \( \square \)

**Claim 16.20**

Claim (16.20) follows from this lemma.

**Lemma 16.8** There exists a universal constant \( c \geq 3 \) such that for all integers \( 1 \leq k \leq \ell \leq \ln(n) \) and \( n \geq c \),
\[
\left( \mathbb{E} \left[ \|R_k^{(\ell)}\|^2 \right] \right)^{1/m} \leq \frac{d^{3/2}}{n} \ln(n)^{23} \ell,
\]
where \( m = \ln(n)/(24 \ln(\ln(n))) \).

The proof is exactly the same than the proof of Lemma 15.4. The extra factor \( \sqrt{d} \) comes from the same reason than in the proof of Lemma 16.5 (the scaling by \( (n/d)^\ell \) while the paths are of length \( \ell + 1 \)). We omit the proof.

**Claim (16.21)**

We have \( |\phi_j^+|_\infty \leq |\phi_j|_\infty / \sqrt{d} \leq b / \sqrt{nd} \). Claim 16.21 follows from this lemma and the union bound.
Lemma 16.9 There exists a universal constant $c \geq 3$ such that for all integers $1 \leq k \leq \ln(n)$, $n \geq cK^{42}$ and vectors $\psi \in \mathbb{R}^{E_{\gamma}}$, 

$$
\left( E \left\| B^{(k)} \psi \right\|^2 \right)^{\frac{1}{2m}} \leq \sqrt{nd} \left\| \psi \right\|_{\infty} \ln(n)^{m} K^{d k},
$$

where $m = \ln(n)/(12 \ln(\ln(n)))$.

Note the lemma improves on Proposition 16.4 and the bound $\left\| B^{(k)} \right\| \cdot \left\| \psi \right\| \leq n \left\| B^{(k)} \right\| \cdot \left\| \psi \right\|_{\infty}$ by a crucial factor $\sqrt{n}$ (up to logarithmic factors).

The proof is again a variation around the proofs of Lemmas 15.1 and 15.4. We write

$$
\left\| B^{(k)} \right\| \cdot \left\| \psi \right\| = \sum_{\gamma} (B^{(k)} \psi)(\gamma)^{2} = \sum_{e, f, f'} (B^{(k)})_{ef} \psi(f)(B^{(k)})_{ef'}\psi(f').
$$

We raise the above expression to a power $m$. The computation leading to (16.23) gives

$$
E \left\| B^{(k)} \psi \right\|^{2m} \leq \left( \frac{n}{d} \right)^{2km} \sum_{\gamma} E \prod_{i=1}^{2m} M_{\gamma_{i}, \gamma_{i+1}} \prod_{t=1}^{k} P_{\gamma_{t}, \gamma_{t+1}} \psi(\gamma_{t+1}),
$$

where the sum is over all $\gamma = (\gamma_{0}, \ldots, \gamma_{2m})$ with $\gamma_{i} = (\gamma_{i, 0}, \ldots, \gamma_{i, k+1}) \in F^{k+1}$ and the boundary conditions: for all $i \in [m]$,

$$
(\gamma_{2i+1, 0}, \gamma_{2i+1, 1}) = (\gamma_{2i, 0}, \gamma_{2i, 1}).
$$

We repeat the definitions in the proof of Lemma 16.5. We associate to an element $\gamma$ as above, an undirected graph $G_{\gamma} = (V_{\gamma}, E_{\gamma})$ with vertices $V_{\gamma} = \{\gamma_{i, t} : 1 \leq i \leq 2m, 0 \leq t \leq k+1\}$ and edge set $E_{\gamma} = \{(\gamma_{i, t}, \gamma_{i, t+1}) : 1 \leq i \leq 2m, 0 \leq t \leq k\}$. From the above boundary conditions, the graph $G_{\gamma}$ has at most $m$ connected components. In particular,

$$
|E_{\gamma}| - |V_{\gamma}| + m \geq 0.
$$

Each edge $e \in E_{\gamma}$ has a multiplicity $m_{e}$ defined as the number of times it is visited by $\gamma$. By construction, (16.24) holds. We again define $\tilde{m}_{e} \leq m_{e}$ as the number of visits of $e \in E$ for some $\{\gamma_{i, t}, \gamma_{i, t+1}\}$ with $t \in [k]$.

The computation leading to 16.25 gives

$$
E \left\| B^{(k)} \psi \right\|^{2m} \leq \sum_{\gamma \in \tilde{W}_{k+1, m}} \left( \frac{n}{d} \right)^{2km - |E_{\gamma}|} \prod_{e \in E_{\gamma}} |P_{e}| \tilde{m}_{e} \cdot \left\| \psi \right\|_{\infty}^{2m},
$$

where $\tilde{W}_{k+1, m}$ is the set of paths $\gamma$ as above such that each edge of $E_{\gamma}$ is visited at least twice.

We define $\tilde{V}_{k+1, m}(s, a)$ as the set of equivalence classes of paths in $\tilde{W}_{k+1, m}$ with $|V_{\gamma}| = s$ and $|E_{\gamma}| = a$. 

 Springer
Lemma 16.10 Let \( a, s \geq 1 \) be integers such that \( a - s + m \geq 0 \). We have
\[
|\tilde{W}_{k,m}(s,a)| \leq (2km)^{6m(a-s+m)+5m}.
\]

Proof The proof is a variant of the proof of Lemma 15.6. Let \( T = \{(i,t) : 1 \leq i \leq 2m, 0 \leq t \leq k-1\} \) ordered with the lexicographic order. For \( \tau = (i,t) \in T \), we set \( e_\tau = \{\gamma_{i,t}, \gamma_{i,t+1}\} \) and \( y_\tau = \gamma_{i,t+1} \). We build the same growing subforest \( F_\tau \) which is a spanning forest of the graph spanned by edges \( (e_s)_{s \leq \tau} \) seen so far.

For each \( i \in [2m] \), we use the same long and short cycling important times. Finally, there are merging times such that two connected components of the graph seen so far merge. Since there are most \( m \) connected components, there are at most \( m - 1 \) merging times. Also, there are most \( g = a - s + m \) excess edges in any connected component of \( G_{y_\tau} \). It follows that, for each \( i \), there are either at most \( g - 1 \) long cycling time and 1 short cycling time, or \( g \) long cycling times and 0 short cycling time.

There are at most \( 2km \) ways to position those important and merging times. There are \( s^2 \) possibilities for the mark of a long cycling time or merging time and at most \( s^2k \) for a short cycling time. We deduce that
\[
|\tilde{W}_{k,m}(s,a)| \leq (2km)^{2mg+m-1(s^2)^2m(g-1)+m-1(s^2_k)^2m}.
\]
We finally use \( s \leq 2km \).

The proof of lemma 16.7 applies verbatim to this case as well.

Lemma 16.11 Let \( y \in \tilde{W}_{k+1,m} \) such that \( |V_y| = s \) and \( |E_y| = a \). We have
\[
\sum_{\gamma' : \gamma' \sim y} \prod_{e \in E_{y'}} |P_e| \tilde{m}_e \leq n^{2km+s} K^{km-a} K^{6(a-s)+8m} \rho^{km}.
\]

We are ready for the proof of Lemma 16.9.

Proof of Lemma 16.9 Note that \( \tilde{W}_{k+1,m}(s,a) \) is empty unless \( 0 \leq s - m \leq a \leq (k+1)m \) (since each edge has multiplicity at least 2 and since 16.29 holds). From 16.30, we get
\[
E|B(k)\psi|^2 \leq \sum_{a=1}^{(k+1)m} \sum_{s=1}^{a+m} \left( \frac{n}{d} \right)^{2km-a} |\tilde{W}_{k+1,m}(s,a)| \max_{y \in \tilde{W}_{k+1,m}(s,a)} \gamma' : \gamma' \sim y \prod_{e \in E_{y'}} |P_e| \tilde{m}_e.
\]
We use Lemmas 16.10, 16.11 and follow the computation of Lemma 15.1. Replacing the summation over \( s \) as a summation over \( g = a - s + m \), we find
\[
E|B(k)\psi|^2 \leq n^m (2km)^{5m} K^{2m} \sum_{a=1}^{(k+1)m} \left( \frac{K}{d} \right)^{km-a} \sum_{g=0}^{\infty} \left( \frac{K^6 (2(k+1)m)^{6m}}{n^g} \right).
\]
The conclusion follows easily.
Claim (16.22)

The proof of Claim 16.22 follows again from the same arguments.

We first define a new matrix \( \tilde{S} \in M_{En}(\mathbb{R}) \) defined for \( e = (x, y) \) and \( f = (a, b) \), by

\[
\tilde{S}_{ef} = S_{ef} M_f = 1_{\{a=x\} \cup \{y=b\}} \frac{P_{ya}}{d} M_{ab}.
\]

Since \( M_e^2 = M_e \), we find

\[
\tilde{B}^{(k-1)} S \Delta B^{(\ell-k-1)} = \tilde{S} \Delta B^{(\ell-k-1)}.
\]

We then write

\[
\| \tilde{B}^{(k-1)} S \Delta B^{(\ell-k-1)} \| \leq \| \tilde{B}^{(k-1)} \| \| \tilde{S} \| \| \Delta B^{(\ell-k-1)} \|.
\]

(16.31)

We may estimate the operators norms on right-hand side of (16.31) separately.

The statement of Lemma 16.5 applies unchanged to \( \tilde{B}^{(k-1)} \). We thus find

\[
\left( \mathbb{E} \left\{ \| \tilde{B}^{(k)} \|^{2m} \right\} \right)^{\frac{1}{2m}} \leq \ln (n) \frac{9 \sqrt{d} K^4 \vartheta^k}{L},
\]

(16.32)

where \( m = \ln (n/K^6)/(12 \ln (\ln (n))) \).

The computation leading to 16.25 applied to \( B^{(k)} \) gives

\[
\mathbb{E} \| \Delta B^{(k)} \|^{2m} \leq d^{2m} \sum_{\gamma \in \tilde{W}_{k+1,m}} \left( \frac{n}{d} \right)^{2(k+1)m-|E_{\gamma}|} \prod_{e \in E_{\gamma}} |P_{e}|^{m_e},
\]

(16.33)

where \( \tilde{W}_{k+1,m} \) is the set of paths \( \gamma \) as below (16.23). The only difference between \( \tilde{W}_{k+1,m} \) and \( \tilde{W}_{k+1,m} \subset \tilde{W}_{k+1,m} \) is that there is no constraint on edge multiplicities. We note that Lemma 16.6 holds also for the equivalence classes of \( \tilde{W}_{k+1,m} \) since the proof does not use the edge multiplicities. Due to the boundary conditions, there at most \( 2km \) edges visited by a path \( \gamma \in \tilde{W}_{k+1,m} \). We also use the rough bound \( |P_{e}| \leq L/n \). We deduce from 16.33 that

\[
\mathbb{E} \| \Delta B^{(k)} \|^{2m} \leq d^{2m} \sum_{a=1}^{2km} \sum_{s=1}^{a+1} \left( \frac{n}{d} \right)^{2(k+1)m-a} n^s (2km)^{6m(a-s+1)+2m} \left( \frac{L}{n} \right)^{2(k+1)m}.
\]

We deduce that

\[
\left( \mathbb{E} \left\{ \| \Delta B^{(k)} \|^{2m} \right\} \right)^{\frac{1}{2m}} \leq \ln (n)^9 L^{k+1},
\]

(16.34)

where \( m = \ln (n)/(12 \ln (\ln (n))) \).
Finally, we have for any \( e = (x, y) \in E_n \),

\[
\sum_{f \in E_n} |(\tilde{S} \tilde{S}^*)_{ef}| \leq \sum_{f, g \in E_n} |S_{eg} S_{fg}| M_g \\
\leq \frac{L^2}{nd^2} (2 \deg(x) + 2 \deg(y)) \leq \frac{4L^2}{nd^2} \max \deg(o),
\]

where \( \deg(x) = \sum_{y} M_{xy} \) is the degree of vertex \( x \in [n] \) in the random graph \( G \). We deduce from 11.7 that with probability \( 1 - 1/n \), for some some universal constant \( c > 0 \), we have for all \( e \in E_n \)

\[
\|\tilde{S}\tilde{S}^*\| \leq \max_{e \in E_n} \sum_{f \in E_n} |(\tilde{S} \tilde{S}^*)_{ef}| \leq \frac{cL^2}{nd} \ln(n).
\]

This implies that

\[
\|S\| \leq \sqrt{\frac{cL^2}{nd} \ln(n)}.
\]

We thus obtain from this last bound and 16.32–16.33–16.34 that for all \( k \in [\ell - 1] \) with probability at least \( 1 - 3/n \), we have, for some \( c > 0 \),

\[
\sqrt{n} \| \tilde{B}^{(k-1)} S \Delta B^{(\ell-k-1)} \| \leq c \ln(n)^{19} K^{4} \delta^{k-1} L^{\ell-k+1} \leq c \ln(n)^{19} K^{4} \delta^{\ell} d^{\ell-k+1}.
\]

where we have used that \( L = \delta_1 d \). We obtain claim (16.22).

17 Proofs for the Rectangular case

17.1 Proof for Lemma 6.1

**Proof** We recall that two Bernoulli random variables \( B_1, B_2 \) are independent with parameters \( b_1, b_2 \) if and only if \( \Pr(B_1 = 1, B_2 = 1) = b_1b_2 \).

For the proof we note \( \delta = d/n \) and \( p = (1 - q)/2 \). Clearly, entries above the diagonal and on the diagonal of \( \tilde{M} \) are independent and their distribution is Bernoulli with parameter \( \delta p + \delta q = q/(p + q) = 2q/(1 + q) \), where we used 6.5.

We must check that the entries \( M_{x,y} \) and \( M_{y,x} \) are themselves independent when \( x \neq y \). We have \( \Pr(\tilde{M}_{x,y} = 1, \tilde{M}_{y,x} = 1) = \delta q \). Consequently, the entries of \( \tilde{M} \) are independent if and only if \( \delta q = (\delta p + \delta q)^2 \), or equivalently if \( q \) satisfies 6.5. \( \square \)
17.2 Link Between a Matrix and its Hermitization

The link between a matrix $P$ and its so-called symmetric dilation
\[
\begin{pmatrix}
0 & P \\
P^* & 0
\end{pmatrix}
\]
is well-known in the literature on random matrices. We let the proofs to the reader—they are mainly verifications.

**Lemma 17.1** (structure of $\tilde{P}$) *Let us write the singular value decomposition of $P$ in the following way:

\[
P = \sum_{i=1}^{\text{rank}(P)} \sigma_i \xi_i \xi_i^*
\]

where

- $\sigma_1 \geq \cdots \geq \sigma_{\text{rank}(P)} > 0$ are the singular values,
- $\xi_i \in \mathbb{C}^m$ is an orthonormal family of left singular vectors,
- $\xi_i \in \mathbb{C}^n$ is an orthonormal family of right singular vectors.

Then a spectral decomposition of the Hermitian matrix $\tilde{P}$ is

\[
\tilde{P} = \sum_{i=1}^{\text{rank}(P)} \sigma_i \varphi_i^+ (\varphi_i^+)^* - \sum_{i=1}^{\text{rank}(P)} \sigma_i \varphi_i^- (\varphi_i^-)^*
\]

(17.1)

where the $\varphi_i^\pm \in \mathbb{C}^{m+n}$ are the orthonormal vectors defined by

\[
\varphi_i^+ = \frac{1}{\sqrt{2}} \left( \xi_i \right), \quad \varphi_i^- = \frac{1}{\sqrt{2}} \left( -\xi_i \right).
\]

(17.2)

17.3 Proof of Theorem 6.8

We keep the notations and setting of the theorem and we place ourselves on the event of Theorem 6.5. Let us decompose the unit right-eigenvectors $\psi_i^+$ associated with the eigenvalue $+\lambda_i$ into their first $m$ components and their last $n$ ones: for $i \in [\tilde{r}_0]$:

\[
\psi_i^+ = \begin{pmatrix} \psi_{i,1}^+ \\ \psi_{i,2}^+ \end{pmatrix}
\]

(17.3)

where $\psi_{i,1} \in \mathbb{C}^m$ and $\psi_{i,2} \in \mathbb{C}^n$ and $|\psi_{i,1}|^2 + |\psi_{i,2}|^2 = |\psi_i|^2 = 1$. We saw in Lemma 6.7 that the eigenvectors $\psi_i^-$ are linked with these ones by

\[
\psi_i^- = \begin{pmatrix} -\psi_{i,1}^- \\ \psi_{i,2}^- \end{pmatrix}.
\]

(17.4)
In the course of the proof, we are going to need an important proposition.

**Proposition 17.2** On the event of Theorem 6.5, we have

\[
\left| \psi_{i,1} \right|^2 - \frac{\gamma_i}{2 \gamma_i} \leq \frac{C_0 \tilde{\tau}_0^\ell}{1 - \tilde{\tau}_i, \ell} \quad \left| \psi_{i,2} \right|^2 - \frac{\gamma_i}{2 \gamma_i} \leq \frac{C_0 \tilde{\tau}_0^\ell}{1 - \tilde{\tau}_i, \ell}.
\] (17.5)

**Proof of Proposition 17.2** The eigenvectors \( \psi_i^\pm \) have unit norms, hence \( \langle \psi_i^+, \psi_i^- \rangle = 1 \), and thanks to Theorem 6.5 we also have \( \langle \psi_i^+, \psi_i^- \rangle \approx \Gamma_i, i^- / \gamma_i \). Both equations can be written

\[
\left| \psi_{i,1} \right|^2 + \left| \psi_{i,2} \right|^2 = 1 \quad - \left| \psi_{i,1} \right|^2 + \left| \psi_{i,2} \right|^2 \approx \frac{\Gamma_{i, i^+}^-}{\gamma_i}.
\]

Adding/subtracting both equations yields

\[
\left| \psi_{i,1} \right|^2 \approx \frac{\gamma_i - \Gamma_{i, i^+}^-}{2 \gamma_i} \quad \left| \psi_{i,2} \right|^2 \approx \frac{\gamma_i + \Gamma_{i, i^+}^-}{2 \gamma_i}.
\] (17.6)

We can now turn to the proof of the theorem.

**Proof of Theorem 6.8** Let us fix \( i, j \in \tilde{r}_0 \). By 6.22, we have \( \langle \psi_i^+, \varphi_j^+ \rangle \approx \delta_{i,j} \sqrt{\gamma_i / \gamma_j} \), and \( \langle \psi_i^+, \varphi_j^- \rangle \approx 0 \), which translates into

\[
\langle \psi_{i,1}, \xi_j \rangle + \langle \psi_{i,2}, \xi_j \rangle \approx \delta_{i,j} \sqrt{2 / \gamma_i} \quad - \langle \psi_{i,1}, \xi_j \rangle + \langle \psi_{i,2}, \xi_j \rangle \approx 0.
\] (17.7)

Adding or subtracting both approximations yields

\[
\langle \psi_{i,1}, \xi_j \rangle \approx \delta_{i,j} \sqrt{\frac{1}{2 \gamma_i}} \quad \langle \psi_{i,2}, \xi_i \rangle \approx \delta_{i,j} \sqrt{\frac{1}{2 \gamma_i}}.
\] (17.8)

Similarly, by 6.23, we have \( \langle \psi_i^+, \psi_j^+ \rangle \approx \Gamma_{i,j}^{++, \pm} / \sqrt{\gamma_i \gamma_j} \) and \( \langle \psi_i^+, \psi_j^- \rangle \approx \Gamma_{i,j}^{++, -} / \sqrt{\gamma_i \gamma_j} \), which translates into

\[
\langle \psi_{i,1}, \psi_{j,1} \rangle + \langle \psi_{i,2}, \psi_{j,2} \rangle \approx \frac{\Gamma_{i,j}^{++, \pm}}{\sqrt{\gamma_i \gamma_j}} \quad - \langle \psi_{i,1}, \psi_{j,1} \rangle + \langle \psi_{i,2}, \psi_{j,2} \rangle \approx \frac{\Gamma_{i,j}^{++, -}}{\sqrt{\gamma_i \gamma_j}}.
\] (17.9)

Here again, we can subtract or add both identities and we obtain

\[
\langle \psi_{i,1}, \psi_{j,1} \rangle \approx \frac{\Gamma_{i,j}^{++, \pm} - \Gamma_{i,j}^{++, -}}{2 \sqrt{\gamma_i \gamma_j}} = \frac{\Gamma_{i,j}^{\Delta}}{2 \sqrt{\gamma_i \gamma_j}}.
\] (17.10)
\[
\langle \psi_{i,2}, \psi_{j,2} \rangle \approx \frac{\Gamma_{i,j}^{-} + \Gamma_{i,j}^{+}}{2\sqrt{\gamma_i \gamma_j}} = \frac{\Gamma_{i,j}^{\nabla}}{2\sqrt{\gamma_i \gamma_j}}.
\] (17.11)

Similarly, by 6.24, we have \(\langle \psi_i^+, \psi_j^+ \rangle \approx \delta_{i,j}/\gamma_i\) and \(\langle \psi_i^+, \psi_j^- \rangle \approx 0\), which translates into
\[
\langle \psi_{i,1}, \psi_{j,1} \rangle + \langle \psi_{i,2}, \psi_{j,2} \rangle \approx \frac{\delta_{i,j}}{\gamma_i} - \langle \psi_{i,1}, \psi_{j,1} \rangle + \langle \psi_{i,2}, \psi_{j,2} \rangle \approx 0.
\]

Adding/subtracting yields
\[
\langle \psi_{i,1}, \psi_{j,1} \rangle \approx \frac{\delta_{i,j}}{2\gamma_i}, \quad \langle \psi_{i,2}, \psi_{j,2} \rangle \approx \frac{\delta_{i,j}}{2\gamma_i}.
\] (17.12)

Thanks to Lemma 6.7 and 17.5, we know that up to some sign,
\[
\chi_i = \frac{\psi_{i,1}}{|\psi_{i,1}|} \approx \psi_{i,1} \sqrt{\frac{2\gamma_i}{\gamma_i^\Delta}}, \quad \pi_i = \frac{\psi_{i,2}}{|\psi_{i,2}|} \approx \psi_{i,2} \sqrt{\frac{2\gamma_i}{\gamma_i^\nabla}}.
\]

Combining this with 17.8–17.10–17.11–17.12 closes the proof of all the claims contained in Theorem 6.8.

\[\square\]

### 17.4 Proof of Theorem 6.10

The first two identities directly follow from 6.32. For the third identity we have
\[
\langle \hat{\xi}_{i,\text{avg}}, \hat{\xi}_{j,\text{avg}} \rangle = \langle \chi_i, \chi_j \rangle + \langle \chi_i', \chi_j \rangle + \langle \chi_i, \chi_j' \rangle + \langle \chi_i', \chi_j' \rangle \approx \sqrt{\frac{1}{\gamma_i^\Delta}} + \sqrt{\frac{1}{\gamma_j^\Delta}} \approx \sqrt{\frac{2}{\gamma_i^\Delta + 1}}
\]
and the proof for\(\hat{\xi}_{i,\text{avg}}\) is the same. We now turn to the correlations between these estimators. Clearly,
\[
\langle \hat{\xi}_{i,\text{sim}}, \hat{\xi}_{j,\text{sim}} \rangle = \langle \chi_i, \chi_j \rangle = \frac{\Gamma_{i,j}^\Delta}{\sqrt{\gamma_i^\Delta \gamma_j^\Delta}}.
\]

For the averaged estimators, we have
\[
\langle \hat{\xi}_{i,\text{avg}}, \hat{\xi}_{j,\text{avg}} \rangle = \frac{\langle \chi_i, \chi_j \rangle + \langle \chi_i', \chi_j \rangle + \langle \chi_i, \chi_j' \rangle + \langle \chi_i', \chi_j' \rangle}{\sqrt{4(1 + \langle \chi_i, \chi_j' \rangle)(1 + \langle \chi_j, \chi_j' \rangle)}}
\]
\[ \approx 2 \frac{\Gamma_{i,j}^\Delta}{\sqrt{n_i n_j}} + 2 \delta_{i,j} \frac{\gamma_i^\Delta}{\gamma_j^\Delta}. \]

When \( i = j \) this is obviously equal to 1, and when \( i \neq j \) it is equal to

\[ \frac{2 \Gamma_{i,j}^\Delta}{\sqrt{n_i n_j}} \frac{1}{(1 + 1/\gamma_i^\Delta)(1 + 1/\gamma_j^\Delta)}. \]

17.5 Proof of Proposition 7.3

We recall the setting: here, \( P \) has rank one and can be written \( P = \zeta \xi^* \) where \( \zeta, \xi \) are unit vectors. We recall that \( \tilde{Q} = (n + m) \begin{pmatrix} 0 & (\zeta \xi^*) \odot (\zeta \xi^*)^* \\ (\zeta \xi^*) \odot (\zeta \xi^*)^* & 0 \end{pmatrix} \).

The SVD of \((\zeta \xi^*) \odot (\zeta \xi^*)\) will be written \( \kappa \tilde{\zeta} \tilde{\xi}^* \) where \( \kappa := |\zeta|^2|\xi|^2 \) and \( \tilde{\zeta}, \tilde{\xi} \) are the unit singular vectors:

\[ \tilde{\zeta} = \frac{\zeta^2}{|\zeta|^2_4}, \quad \tilde{\xi} = \frac{\xi^2}{|\xi|^2_4}. \]

where we used \( \zeta^2 \) for the entry-wise product \( \zeta \odot \zeta \), same thing for \( \xi^2 \). The operator norm \( \tilde{\rho} = \| \tilde{Q} \| \) is equal to

\[ \tilde{\rho} = (n + m)\kappa = n(1 + \alpha)|\zeta|^2_4|\xi|^2_4. \quad (17.13) \]

The spectral decomposition of \( \tilde{Q} \) is

\[ \tilde{Q} = \tilde{\rho} \phi^+ (\phi^+)^* - \tilde{\rho} \phi^- (\phi^-)^*, \quad (17.14) \]

where \( \phi^\pm \) are the unit eigenvectors of \( \tilde{Q} \) which are given by

\[ \phi^\pm = \frac{1}{\sqrt{2}} \begin{pmatrix} \tilde{\zeta} \\ \pm \tilde{\xi} \end{pmatrix}. \]

In particular, the powers of \( \tilde{Q} \) are given by \( \tilde{Q}^s = \tilde{\rho}^s \phi^+ (\phi^+)^* + (-1)^s \tilde{\rho}^s \phi^- (\phi^-)^* \).

Our detection threshold is

\[ \max\{\tilde{\delta}_1, \tilde{\delta}_2\} = \tilde{\delta}_1 = \sqrt{\frac{\tilde{\rho}}{d}} = \sqrt{\frac{2n(1 + \alpha)(|\xi|^4_4|\xi|^4_4)^2}{(1 + \alpha)d}} = \sqrt{\frac{2n|\zeta|^2_4|\xi|^2_4}{d}}. \]
We now compute \( \gamma_i^\Delta = \Gamma_{i,i}^\Delta \), which is defined in 6.29 by

\[
\gamma_i^\Delta = \sum_{s=0}^\ell (\tilde{Q}^\Delta_i, \zeta_i \Delta \zeta_i) \frac{d^s}{ds}.
\]  

To do the computation in 17.15 we are going to need a few steps. First, we have

\[
\langle \phi^\pm, \mathbf{1} \rangle = |\zeta|_4^2 \pm |\xi|_4^2 \sqrt{2} |\xi|_4^2 =: \Upsilon_\pm.
\]

Second, we have

\[
\langle \phi^\pm, \zeta \Delta \zeta \rangle = \frac{1}{\sqrt{2}} \left( \pm \langle \zeta, \zeta^2 \rangle \right)
= \pm |\xi|_4^2 \sqrt{2} =: \Xi_\pm.
\]

Third, using the spectral decomposition of \( \tilde{Q} \), we have

\[
\langle \tilde{Q}^\Delta_i, \zeta \Delta \zeta \rangle = \tilde{\rho}^s \langle \phi^+, \zeta \Delta \zeta \rangle \langle \phi^+, \mathbf{1} \rangle + (-1)^s \tilde{\rho}^s \langle \phi^-, \zeta \Delta \zeta \rangle \langle \phi^-, \mathbf{1} \rangle
= \tilde{\rho}^s \Xi_+ \Upsilon_+ + (-1)^s \tilde{\rho}^s \Xi_- \Upsilon_-.
\]

When we sum over \( s \) in 17.15 and we gather the preceding identities, we obtain

\[
\gamma_1^\Delta = \Xi_+ \Upsilon_+ \sum_{s=0}^\ell \frac{\tilde{\rho}^s}{d^s} + \Xi_- \Upsilon_- \sum_{s=0}^\ell \frac{(-1)^s \tilde{\rho}^s}{d^s}
= \Xi_+ \Upsilon_+ \sum_{s=0}^\ell \tilde{\Theta}_1^2 s + \Xi_- \Upsilon_- \sum_{s=0}^\ell (-1)^s \tilde{\Theta}_1^2 s
= \Xi_+ \Upsilon_+ \frac{1 - \tilde{\Theta}_1^2 (\ell + 1)}{1 - \tilde{\Theta}_1^2} + \Xi_- \Upsilon_- \frac{1 - (-1)^\ell \tilde{\Theta}_1^2 (\ell + 1)}{1 - \tilde{\Theta}_1^2}.
\]

When computing the asymptotic value of \( \gamma_+ \) as \( n \) is large, we can neglect \( \tilde{\Theta}_1^2 \). Moreover, simple manipulations show that

\[
\Xi_+ \Upsilon_+ = \frac{|\xi|_4^2 + |\xi|_4^2}{2 |\xi|_4^2}, \quad \Xi_- \Upsilon_- = \frac{|\xi|_4^2 - |\xi|_4^2}{2 |\xi|_4^2}.
\]

We find

\[
\gamma_i^\Delta \sim \frac{\Xi_+ \Upsilon_+}{1 - \tilde{\Theta}_1^2} + \frac{\Xi_- \Upsilon_-}{1 + \tilde{\Theta}_1^2}.
\]
\[
\frac{1}{2(1 - \tilde{\vartheta}_1^2)} \left\| \xi \right\|^2_4 + \frac{1}{2(1 + \tilde{\vartheta}_1^2)} \left\| \xi \right\|^2_4 = \frac{1}{2} \left( \left\| \xi \right\|^2_4 + \left\| \xi \right\|^2_4 \right)
\]

as requested in (7.9). The other identity (7.10) is proved in the same way.

### 17.6 Proof of Proposition 2.2 on Mean Square Errors

**Proof** For the sake of this proof only, we will note
\[
P_0 = \frac{U}{\Delta_1} V^* \quad \text{and} \quad \hat{P} = \frac{\hat{U}}{\Delta_1} \hat{V}^*.
\]
where the columns of \(U\) (resp \(\hat{U}\)) are \(\zeta_i\) (resp \(\hat{\zeta}^#_i\)), where \(\Delta = \text{diag}(\sigma_i)\) and \(\hat{\Delta} = \text{diag}(w^#_i)\).

We have

\[
\text{MSE} = \left\| P_0 - \hat{P} \right\|_F^2 = \text{tr}\left((U \Delta V^* - \hat{U} \hat{\Delta} \hat{V}^*)(V \Delta U^* - \hat{V} \hat{\Delta} \hat{U}^*)\right)
\]

By unitary invariance, the first of these four terms is equal to \(\text{tr}(\Delta^2) = \sum \sigma_i^2\). For the other terms, we will use the approximations of Theorem 6.10. We will abuse the symbol \(\approx\) to get rid of the error terms (they can rigorously be neglected). Thanks to Theorem 6.10, we know that \(V^* \hat{V} \approx \nabla\) and \(U^* \hat{U} \approx \nabla'\) where \(\nabla, \nabla'\) are the diagonal matrices whose \(i\)-th entry is equal to \(c^#_{1,i}\) or \(c^#_{2,i}\).

Consequently, the second and third terms in the trace above are close to

\[
-\text{tr}(\Delta \nabla \hat{\Delta} \nabla') \approx -\sum_{i=1}^{\tilde{r}_0} \sigma_i w_i \sigma_i c^#_{1,i} c^#_{2,i} \approx -\sum_{i=1}^{\tilde{r}_0} (\sigma_i c^#_{1,i} c^#_{2,i})^2.
\]

Moreover, from the second part of Theorem 6.10 we have

\[
\hat{V}^* \hat{V} \approx \mathcal{C}_2^# \quad \text{and} \quad \hat{U}^* \hat{U} \approx \mathcal{C}_1^#,
\]

consequently \(\text{tr}(\hat{\Delta} \hat{V}^* \hat{V} \hat{\Delta} \hat{U}^*) \approx \text{tr}(\hat{\Delta} \mathcal{C}_2 \hat{\Delta} \mathcal{C}_1)\). This can be written as

\[
\sum_{i,j} \sigma_i \sigma_j c^#_{1,i} c^#_{2,j}(\mathcal{C}_1)_{i,j}(\mathcal{C}_2)_{i,j}.
\]

As requested. \(\square\)
Some Possible Improvement of the Quantitative Bounds

18.1 On the Threshold $\vartheta_1$

We have already discussed that the threshold $\vartheta_2$ is an intrinsic limitation of the problem as it is the Kesten-Stigum bound. The threshold $\vartheta_1$ however is more artificial and, in most cases, it can be reduced to a smaller value $\vartheta_1' = L'/d$ with $L' \leq L$ at the cost of more technicalities. In this section, we discuss a few ways to improve on this new parameter $L'$ by some recipes. The way to formalize them would depend on the structure of the matrix $P$. We discuss the case of the square matrix $A$ but the same comments extend to the non-backtracking matrix $B$.

A first possibility is the following: we consider the deterministic set $S \subset [n]^2$ of the $o(n/d)$ largest entries of the matrix $P$ in modulus. Then with high probability, the revealed entries (non-zeros entries of $M$) will not intersect $S$. In particular, on this good event, the conclusions of Theorem 4.3 remain valid if we replace $L$ by $L' = \max_{(x,y) \notin S} n |P_{x,y}|$.

In the proofs, we have many times bounded $P_{x,y}$ by $L/n$ and $Q_{x,y}$ by $K/(n\rho)$. A finer probabilistic and combinatorial analysis could be performed by partitioning the entries of $P$ into two or more sets depending on the value of $n|P_{xy}|$. An analysis of this kind was done in [15] (see the parameter $\delta$ there).

Another possibility is to define, for a given $L' \leq L$, the matrix $P'$ obtained from $P$ by putting to 0 all entries of $P$ larger than $nL'$. We could then apply Theorem 4.3 directly to $P'$ and then use spectral perturbation theory (such as Hoffman-Wielandt inequality) to guarantee that the spectra of $P$ and $P'$ are close provided that $L - L'$ is not too large.

18.2 On the Parameter $C_0$

The conclusions of Theorems 4.3 and 5.4 are controlled by a parameter $C_0$ and $\tilde{C}_0$. Again, depending on the structure of $P$ there are ways reduce the value of this constant at the cost of some extra technicalities.

The same remarks as in Sect. 18.1 applies. Every time that we have bounded $nP_{xy}$ by $L/n$ and $Q_{x,y}$ by $K/(n\rho)$, there is room for improvement.

Finally, the incoherence parameter $b$ is important for the eigenvectors $i \in [r_0]$. Since we are interested in an average quantity such as the scalar product $\langle \psi_i, \phi_i \rangle$, it is possible to relax the notion of incoherence by partitioning the entries $x \in [n]$ depending on the value of $\sqrt{n}|\phi_i(x)|$. This makes the application of the general concentration bound Theorem 12.5 more tedious. For the other eigenvectors, $i \notin [n]\setminus[r_0]$, their importance is weighted by their corresponding eigenvalue and there is also room for improvement here.

Acknowledgements CB was supported by ANR-16-CE40-0024-01. SC is supported by ERC NEMO, under the European Union’s Horizon 2020 research and innovation programme Grant Agreement Number 788851. RRN’s work was supported by ONR Grant N00014-15-1-2141, DARPA Young Faculty Award D14AP00086, and ARO MURI W911NF-11-1-039CB. CB and SC thank the University of Michigan for its hospitality in June 2016 and June 2018 where this work was initiated and continued. We thank Walid Hachem.
for his comments. CB and RRN thank Literati Coffee in Ann Arbor, MI for the stimulating environment in which the problem considered here was first brewed.

References

1. D. Achlioptas and F. McSherry, Fast computation of low-rank matrix approximations, Journal of the ACM (JACM), 54 (2007), p. 9.
2. J. Alt, R. Ducatez, and A. Knowles, Extremal eigenvalues of critical Erdős-Rényi graphs, arXiv preprint arXiv:1905.03243 (2019).
3. N. Anantharaman, Some relations between the spectra of simple and non-backtracking random walks, 1703.03852, 2017.
4. O. Angel, J. Friedman, and S. Hoory, The non-backtracking spectrum of the universal cover of a graph, Transactions of the American Mathematical Society, 367 (2015), pp. 4287–4318.
5. S. Arora, Y. Li, Y. Liang, T. Ma, and A. Risteski, A latent variable model approach to pmi-based word embeddings, Transactions of the Association for Computational Linguistics, 4 (2016), pp. 385–399.
6. J. Baik, G. Ben Arous, and S. Péché, Phase transition of the largest eigenvalue for nonnull complex sample covariance matrices, The Annals of Probability, 33 (2005), pp. 1643–1697.
7. A. Basak and M. Rudelson, The circular law for sparse non-hermitian matrices, The Annals of Probability, 47 (2019), pp. 2359–2416.
8. F. Bauer and C. Fike, Norms and exclusion theorems., Numerische Mathematik, 2 (1960), pp. 137–141.
9. F. Benaych-Georges, C. Bordenave, and A. Knowles, Spectral radii of sparse random matrices, arXiv e-prints (2017), arXiv:1704.02945.
10. F. Benaych-Georges, C. Bordenave, and A. Knowles, Largest eigenvalues of sparse inhomogeneous Erdős-Rényi graphs, Ann. Probab., 47 (2019), pp. 1653–1676.
11. F. Benaych-Georges and R. Nadakuditi, The eigenvalues and eigenvectors of finite, low rank perturbations of large random matrices, Advances in Mathematics, 227 (2011), pp. 494–521.
12. C. Bordenave, A new proof of Friedman’s second eigenvalue theorem and its extension to random lifts, Ann. Sci. Éc. Norm. Supér., (to appear).
13. C. Bordenave and D. Chafai, Around the circular law, Probability surveys, 9 (2012).
14. C. Bordenave, M. Lelarge, and L. Massoulié, Nonbacktracking spectrum of random graphs: Community detection and nonregular ramanujan graphs, Annals of probability: An official journal of the Institute of Mathematical Statistics, 46 (2018), pp. 1–71.
15. C. Bordenave, Y. Qiu, and Y. Zhang, Spectral gap of sparse bistochastic matrices with exchangeable rows with application to shuffle-and-fold maps. arXiv:1805.06205.
16. C. Bordenave, Y. Qiu, and Y. Zhang, Spectral gap of sparse bistochastic matrices with exchangeable rows with application to shuffle-and-fold maps, ArXiv e-prints, (2018).
17. S. Boucheron, G. Lugosi, and P. Massart, Concentration inequalities, Oxford University Press, Oxford, 2013. A nonasymptotic theory of independence, With a foreword by Michel Ledoux.
18. G. Brito, I. Dumitriu, and K. D. Harris, Spectral gap in random bipartite biregular graphs and its applications, ArXiv e-prints, (2018).
19. E. J. Candès and B. Recht, Exact matrix completion via convex optimization, Found. Comput. Math., 9 (2009), pp. 717–772.
20. E. J. Candès and T. Tao, The power of convex relaxation: near-optimal matrix completion, IEEE Trans. Inform. Theory, 56 (2010), pp. 2053–2080.
21. A. Chakrabarty, S. Chakrabarty, and R. Hazra, Eigenvalues outside the bulk of inhomogeneous Erdős-Rényi random graphs, arXiv preprint arXiv:1911.08244 (2019).
22. S. Chatterjee, Matrix estimation by universal singular value thresholding, Ann. Statist., 43 (2015), pp. 177–214.
23. Y. Chen, C. Cheng, and J. Fan, Asymmetry helps: Eigenvalue and eigenvector analyses of asymmetrically perturbed low-rank matrices, arXiv preprint arXiv:1811.12804 (2018).
24. S. Coste, The Spectral Gap of Sparse Random Digraphs, ArXiv e-prints, (2017).
25. S. Coste, Grandes valeurs propres de graphes aléatoires dilués, PhD thesis, 2019.
26. M. Davenport and J. Romberg, An overview of low-rank matrix recovery from incomplete observations, IEEE Journal of Selected Topics in Signal Processing, 10 (2016), pp. 608–622.
27. P. Drineas and A. Zouzias, A note on element-wise matrix sparsification via a matrix-valued Bernstein inequality, Inform. Process. Lett., 111 (2011), pp. 385–389.
28. J. Eldridge, M. Belkin, and Y. Wang, Unperturbed: spectral analysis beyond Davis-Kahan, arXiv e-prints, (2017), arXiv:1706.06516.
29. U. Feige and E. Ofek, Spectral techniques applied to sparse random graphs, Random Structures & Algorithms, 27 (2005), pp. 251–275.
30. J. Friedman, J. Kahn, and E. Szeméry, On the second eigenvalue of random regular graphs, in Proceedings of the twenty-first annual ACM symposium on Theory of computing, ACM, 1989, pp. 587–598.
31. Z. Füredi and J. Komlós, The eigenvalues of random symmetric matrices, Combinatorica, 1 (1981), pp. 233–241.
32. D. Gamarnik, Q. Li, and H. Zhang, Matrix completion from O(n) samples in linear time, arXiv preprint arXiv:1702.02267 (2017).
33. L. C. García del Molino, K. Pakdaman, and J. Touboul, Real eigenvalues of non-symmetric random matrices: Transitions and Universality, arXiv e-prints, (2016), arXiv:1605.00623.
34. L. Grasedyck, D. Kressner, and C. Tobler, A literature survey of low-rank tensor approximation techniques, GAMM-Mitteilungen, 36 (2013), pp. 53–78.
35. B. Huang, C. Mu, D. Goldfarb, and J. Wright, Provable models for robust low-rank tensor completion, Pacific Journal of Optimization, 11 (2015), pp. 339–364.
36. R. Kannan and T. Theobald, Games of fixed rank: A hierarchy of bimatrix games, Economic Theory, 42 (2010), pp. 157–173.
37. R. H. Keshavan, A. Montanari, and S. Oh, Matrix completion from a few entries, IEEE transactions on information theory, 56 (2010), pp. 2980–2998.
38. M. Krivelevich and B. Sudakov, The largest eigenvalue of sparse random graphs, Combinatorics, Probability and Computing, 12 (2003), pp. 61–72.
39. F. Krzakala, C. Moore, E. Mossel, J. Neeman, A. Sly, L. Zdeborová, and P. Zhang, Spectral redemption in clustering sparse networks, Proceedings of the National Academy of Sciences, 110 (2013), pp. 20935–20940.
40. A. Kundu and P. Drineas, A Note on Randomized Element-wise Matrix Sparsification, ArXiv e-prints, (2014).
41. M. Lelarge and L. Miolane, Fundamental limits of symmetric low-rank matrix estimation, Probability Theory and Related Fields, (2017), pp. 1–71.
42. D. Levin, Y. Peres, and E. Wilmer, Markov chains and mixing times, Providence, R.I. American Mathematical Society, 2009. With a chapter on coupling from the past by James G. Propp and David B. Wilson.
43. L. Massoulié, Community detection thresholds and the weak ramanujan property, in Proceedings of the forty-sixth annual ACM symposium on Theory of computing, 2014, pp. 694–703.
44. F. Metz, I. Neri, and T. Rogers, Spectra of sparse non-hermitian random matrices, arXiv preprint arXiv:1811.10416 (2018).
45. L. Miotlante, Fundamental limits of low-rank matrix estimation: the non-symmetric case, arXiv e-prints, (2017), arXiv:1702.00473.
46. C. Moore, The computer science and physics of community detection: landscapes, phase transitions, and hardness, Bull. Eur. Assoc. Theor. Comput. Sci. EATCS, (2017), pp. 26–61.
47. E. Mossel, J. Neeman, and A. Sly, Stochastic block models and reconstruction, arXiv preprint arXiv:1202.1499 (2012).
48. R. Nadakuditi, Optshrink: An algorithm for improved low-rank signal matrix denoising by optimal, data-driven singular value shrinkage, IEEE Transactions on Information Theory, 60 (2014), pp. 3002–3018.
49. S. O’Rourke, V. Vu, and K. Wang, Random perturbation and matrix sparsification and completion, ArXiv e-prints, (2018).
50. S. Pal and Y. Zhu, Community Detection in the Sparse Hypergraph Stochastic Block Model, arXiv e-prints, (2019), arXiv:1904.05981.
51. A. Saade, F. Krzakala, and L. Zdeborová, Spectral clustering of graphs with the bethe hessian, in Proceedings of the 27th International Conference on Neural Information Processing Systems - Volume 1, NIPS’14, Cambridge, MA, USA, 2014, MIT Press, p. 406-414.
52. A. Saade, F. Krzakala, and L. Zdeborová, Matrix completion from fewer entries: Spectral detectability and rank estimation, in Proceedings of the 28th International Conference on Neural Information Processing Systems - Volume 1, NIPS’15, Cambridge, MA, USA, 2015, MIT Press, p. 1261-1269.
53. N. Stein, A. Ozdaglar, and P. Parrilo, Separable and low-rank continuous games, International Journal of Game Theory, 37 (2008), pp. 475–504.
54. L. Stephan and L. Massoulié, Non-backtracking spectra of weighted inhomogeneous random graphs, 2020.
55. G. W. Stewart and J. Sun, Matrix perturbation theory, Computer Science and Scientific Computing, Academic Press, Inc., Boston, MA, 1990.
56. K. Tikhomirov and P. Youssef, Outliers in spectrum of sparse Wigner matrices, arXiv e-prints, (2019), arXiv:1904.07985.
57. T. Tsiligkaridis and A. O. Hero, Covariance estimation in high dimensions via kronecker product expansions, IEEE Transactions on Signal Processing, 61 (2013), pp. 5347–5360.
58. C. F. Van Loan, The ubiquitous kronecker product, Journal of computational and applied mathematics, 123 (2000), pp. 85–100.
59. C. F. Van Loan and N. Pitsianis, Approximation with kronecker products, in Linear algebra for large scale and real-time applications, Springer, 1993, pp. 293–314.
60. Y. Watanabe and K. Fukumizu, Graph zeta function in the bethe free energy and loopy belief propagation, Advances in Neural Information Processing Systems 22 - Proceedings of the 2009 Conference, 22 (2010).
61. Y. Yu, T. Wang, and R. J. Samworth, A useful variant of the davis–kahan theorem for statisticians, Biometrika, 102 (2014), pp. 315–323.

Publisher's Note Springer Nature remains neutral with regard to jurisdictional claims in published maps and institutional affiliations.