A connected graph is called Hamilton-connected if there exists a Hamiltonian path between any pair of its vertices. Determining whether a graph is Hamilton-connected is an NP-complete problem. Hamiltonian and Hamilton-connected graphs have diverse applications in computer science and electrical engineering. The detour index of a graph is defined to be the sum of lengths of detours between all unordered pairs of vertices. The detour index has diverse applications in chemistry. Computing the detour index for a graph is also an NP-complete problem. In this paper, we study the Hamilton-connectivity of convex polytopes. We construct three infinite families of convex polytopes and show that they are Hamilton-connected. An infinite family of non-Hamilton-connected convex polytopes is also constructed, which, in turn, shows that not all convex polytopes are Hamilton-connected. By using Hamilton connectivity of these families of graphs, we compute exact analytical formulas of their detour index.

1. Introduction and Preliminaries

All graphs in this paper are simple, loopless, finite, and connected.

A cycle in a graph G is called Hamiltonian if it travels all the vertices of G. Moreover, a path in G is called Hamiltonian path if it traverses all the vertices of G. Not every graph contains a Hamiltonian cycle. For instance, any tree is an acyclic graph, so it cannot contain a Hamiltonian cycle. However, a tree may still contain a Hamiltonian path. A graph G is called Hamiltonian if there exist a Hamiltonian cycle in it. By definition, any cycle graph or a clique graph are Hamiltonian. Furthermore, G is called traceable if it contains a Hamiltonian path. Of course all Hamiltonian graphs are traceable. However, there exist graphs which are traceable but not Hamiltonian. The first nontrivial example which comes in mind immediately is the so-called Petersen graph which is traceable but not Hamiltonian.

A graph which contains a Hamiltonian path between every two vertices of G is called Hamilton connected. They were introduced by Ore [1] in 1963. Frucht [2] studied a canonical representation of trivalent Hamiltonian graphs. A bipartite can never be a Hamilton-connected, since there cannot exist a Hamiltonian path between vertices of the same partite set. In that case, a bipartite graph is called Hamilton-laceable if there exist Hamiltonian paths between vertices of different partite sets. There is an extensive literature available on Hamiltonianicity and Hamilton-connectivity of graphs, see, for instance, [3–7].

Chartrand et al. [8] showed that the square of a block graph is Hamilton-connected. Thomasson [9] studied Hamilton-connected tournaments in graphs. Chang et al. [10] studied panconnectivity, fault-tolerant Hamiltonicity, and Hamiltonian-connectivity in alternating group graphs by considering them as interconnection networks. Kewen et al. [11] derived a sufficient condition for a graph to be Hamilton-connected. Zhou and Wang [12] proved certain sufficient conditions for a graph to be Hamilton-connected in terms of the edge number, the spectral radius, and the signless Laplacian spectral radius of the graph. Zhou et al.
[13] calculated the Wiener and Harary indices of Hamilton-connected graphs with large diameter. Wei et al. [14] derived some spectral analogues of Erdős’s theorems for Hamilton-connected graphs. Hung et al. [15] studied Hamilton-connectivity of alphabet grid graphs. Zhou et al. [16] extended a result of Fiedler and Nikiforov and derived signless Laplacian spectral conditions for Hamilton-connected graphs with large minimum degree. Recently, Shabbir et al. [17] studied Hamilton-connectivity in Teoplitz graphs.

By preserving the vertex-edge incidence relation in convex polytopes, their graphs are constructed. Baća [18–20] was the first researcher to consider these families of geometric graphs. In [20] (resp. [19]), Baća studied the problem of magic (resp. graceful and antigraceful) labeling of convex polytopes, whereas, in [18], the problem of face antimagic labeling of convex polytopes was studied. Miller et al. [21] studied the vertex-magic total labeling of convex polytopes. Imran et al. [22–24] computed the minimum metric dimension of various infinite families of convex polytopes. In particular, they showed that these infinite families of convex polytopes have constant metric dimension. Malik et al. [25] also constructed two infinite families of convex polytopes having constant metric dimension. Other closely related infinite families of graphs with constant metric dimension are studied in [26]. Kratica et al. [27] studied the strong metric dimension of certain infinite families of convex polytopes by constructing their doubly resolving sets. The fault-tolerant metric dimension (resp. mixed metric dimension) of convex polytopes was studied by Raza et al. [28] (resp. Raza et al. [29]). Metric dimension of other related families of graphs such as Cayley graph is studied by Vetrik and Abas [30, 31]. The binary locating-dominating number of convex polytopes is studied by Simić et al. [32] and Raza et al. [33]. The open-locating-dominating number of certain convex polytopes has recently been studied by Savic et al. [34]. Liu et al. [35–39] studied application of various graph parameters in electrical networks and related systems. For a graph $G$, let $\ell(x, y)$ be the length of a longest path (i.e., detour) between vertices $x$ and $y$ of $G$. The detour index [40] is defined to be the sum of lengths of all detours between unordered pairs of vertices in $G$. The detour index of a graph $G$ is usually denoted by $\omega(G)$:

$$\omega(G) = \sum_{\{x,y\}\subseteq V(G)} \ell(x, y). \quad (1)$$

The detour index has important applications in chemistry. Applications of this parameter in quantitative structure activity and property relationship models were put forward by Lukovits [41]. Besides giving further applications of the detour index, Trinajstić et al. [42] conducted a comparative analysis with the Wiener index in terms of applicability in correlating structure boiling point of organic compounds. Moreover, its further applications in predicting the normal boiling points of cyclic and acyclic alkanes were studied by Rücker and Rücker [43].

An algorithm for tracing the detour between two vertices in a graph was proposed by Lukovits and Razinger [44]. They also applied their algorithm in detecting detours and computing the detour index of graphs corresponding to fused bicyclic skeletons. Rücker and Rücker [43] and Trinajstić et al. [45] proposed certain computer methods to trace out detours and thus calculation of the detour index of graphs. It has already been shown in [46] that the problem of finding the detour index of a given graph is computationally NP-complete. Trinajstić et al. [45] also proposed a method of calculating the detour matrix of reasonably small sizes. Note that the detour index is equal to the sum of all the entries of the detour matrix dividing by two.

Mahmiani et al. [47] proposed the edge versions of the detour index and studied their mathematical properties. Zhou and Cai [48] proved some upper and lower bounds on the detour index of graphs. Qi and Zhou [49] studied minimum unicyclic graphs with respect to the detour index. Du [50] studied the minimum detour index of bicyclic graphs. Fang et al. [51] characterized the minimum detour index of some families of tricyclic graphs. Karbasioun et al. [52] studied the applications of the detour index in infinite families of nanostar dendrimers. Wu and Deng [53] computed the detour index for a chain of C20 fullerenes. Kaladevi [54] studied spectral properties of the detour index in relation with the Laplacian energy of graphs. Recently, Abdullah and Omar [55] introduced the restricted edge version of the detour index and studied it for some families of graphs.

We end this section with an important and well-known result bounding the detour index in terms of graph parameters.

**Theorem 1** (see [56]). Let $G$ be a connected graph with $n \geq 3$ vertices. Then,

$$\omega(G) \leq \frac{n(n-1)^2}{2}, \quad (2)$$

with left equality if and only if $G \cong S_n$, and with right equality if and only if $G$ is Hamilton-connected.

In this paper, we study the Hamilton connectivity of certain infinite families of convex polytopes. More precisely, we construct three infinite families of Hamilton-connected convex polytopes. Moreover, we construct an infinite family non-Hamilton-connected convex polytope. By using Hamilton-connectivity of these families of graphs, we compute exact formulas of their detour index.

### 2. Hamilton Connectivity and the Detour Index of $R_n$

In this section, we show that the $n$-dimensional infinite family of convex polytopes $R_n$ is Hamilton-connected. Afterwards, we use its Hamilton connectivity to find analytical exact expression of the detour index of the graph $R_n$.

For $n \geq 5$, the graph of convex polytope $R_n$ is defined in [20]. It has the vertex set

$$V(R_n) = \left\{x_j; y_j; z_j; 1 \leq j \leq n\right\}, \quad (3)$$

and edge set
Complexity

\[
E(R_n) = \{(x_j, x_{j+1}):(x_j, y_j); (x_j, y_{j+1}); (y_j, y_{j+1});
\]

\[
(y_j, z_j); (z_j, z_{j+1});: 1 \leq j \leq n\}. \tag{4}
\]

By convention, we suppose that \(x_{n+1} = x_n\), \(y_{n+1} = y_n\), and \(z_{n+1} = z_n\). See Figure 1 to view the \(n\)-dimensional convex polytope \(R_n\).

For a positive integer \(v \in \mathbb{Z}^+\), we write \(v|2\) (resp. \(v|2\)) if \(v\) is even (resp. odd).

**Theorem 2.** The \(n\)-dimensional convex polytope \(R_n\), with \(n \geq 4\), is Hamilton-connected.

\[
P_{H}(u, v): u = z_1 \sim \{z_{n-j}; 0 \leq j \leq n-i-1\} \sim \{y_{i-j}; 0 \leq j \leq i-2\} \sim \{x_j; 3 \leq j \leq i+1\} \sim \{x, y_j; i+2 \leq j \leq n\} \sim x_1y_1x_2y_2
\]

\[
\sim \{z_j; 2 \leq j \leq i\} = v. \tag{5}
\]

**Proof.** We prove this result by definition. For this, we have to show that there exists a Hamiltonian path between any pair of vertices of \(R_n\).

Let \(P_{H}(u, v)\) be a Hamiltonian path between vertices \(u\) and \(v\) in \(R_n\). Let \(V(R_n) = Z \cup Y \cup X\) such that \(Z = \{z_1, z_2, \ldots, z_n\}\), \(Y = \{y_1, y_2 \ldots y_n\}\), and \(X = \{x_1, x_2, \ldots, x_n\}\), see, Figure 1.

Case 1: \(u = z_1\) and \(v = z_i\), \(2 \leq i \leq n\).

Case 1.1: \(2 \leq i \leq n-1\):

\[
P_{H}(u, v): u = z_1 \sim \{z_{n-j}; 0 \leq j \leq n-i-1\} \sim \{y_{i-j}; 0 \leq j \leq n-i\} \sim \{x_j; i+1 \leq j \leq n\} \sim x_1y_1x_2y_2
\]

\[
\sim \{z_j; 2 \leq j \leq i\} = v. \tag{6}
\]

Case 2: \(u = z_1\) and \(v = y_j\), \(1 \leq i \leq n\).

Case 2.1: \(1 \leq i \leq n-1\):

\[
P_{H}(u, v): u = z_1y_1 \sim \{x_j; 1 \leq j \leq n\} \sim \{y_{n-j}; 0 \leq j \leq n-2\}
\]

\[
\sim \{z_j; 2 \leq j \leq n\} \sim y_n = v. \tag{7}
\]

Case 2.2: \(i = n\):

\[
P_{H}(u, v): u = z_1y_1 \sim \{x_j; 1 \leq j \leq n\} \sim \{y_{n-j}; 1 \leq j \leq n-2\}
\]

\[
\sim \{z_j; 2 \leq j \leq n\} \sim y_n = v. \tag{8}
\]

Case 3: \(u = z_1\) and \(v = x_j\), \(1 \leq i \leq n\).

Case 3.1: \(i = 1\):

\[
P_{H}(u, v): u = z_1 \sim \{z_j; 2 \leq j \leq n\} \sim \{y_{n-j}; 0 \leq j \leq n-1\}
\]

\[
\sim \{x_j; 2 \leq j \leq n\} \sim x_1 = v. \tag{9}
\]

Case 3.2: \(2 \leq i \leq n\):

\[
P_{H}(u, v): u = z_1 \sim \{z_j; 2 \leq j \leq n\} \sim \{y_{n-j}; 0 \leq j \leq n-1\}
\]

\[
\sim \{y_{i-j}y_{i-j}^1; 1 \leq j \leq i-1\} \sim \{x_{i-j}; 0 \leq j \leq n-1\} = v. \tag{10}
\]

Case 4: \(u = y_1\) and \(v = z_i\), \(1 \leq i \leq n\).

Case 4.1: \(i = 1\):

\[
P_{H}(u, v): u = y_1 \sim \{x_j; 1 \leq j \leq n\} \sim \{y_{n-j}; 0 \leq j \leq n-2\}
\]

\[
\sim \{z_j; 2 \leq j \leq n\} \sim z_1 = v. \tag{11}
\]

Case 4.2: \(2 \leq i \leq n-1\):

\[
P_{H}(u, v): u = y_1x_1 \sim \{x_j; 2 \leq j \leq i\} \sim \{x_j; i+1 \leq j \leq n\}
\]

\[
\sim \{y_{n-j}; 0 \leq j \leq n-i-1\} \sim \{z_j; i+1 \leq j \leq n\}
\]

\[
\sim \{z_j; 1 \leq j \leq i\} = v. \tag{12}
\]

Case 4.3: \(i = n\):

\[
P_{H}(u, v): u = y_1 \sim \{z_j; 1 \leq j \leq n-1\} \sim \{y_{n-j}; 1 \leq j \leq n-2\}
\]

\[
\sim \{x_j; 2 \leq j \leq n\} \sim x_1y_1x_2y_2
\]

\[
\sim \{z_j; 1 \leq j \leq i\} = v. \tag{13}
\]
Case 5: \( u = y_1 \) and \( v = y_i, \ 2 \leq i \leq n \).

Case 5.1: \( i = 2 \):

\[
P_H(u, v); u = y_1 \sim \{x_j; 1 \leq j \leq n\} \sim \{y_{n-j}; 0 \leq j \leq n - 3\}
\sim \{z_j; 3 \leq j \leq n\} \sim z_1 z_2 y_2 = v.
\]

(14)

Case 5.2: \( 3 \leq i \leq n - 1 \):

\[
P_H(u, v); u = y_1 x_1 \sim \{x_j y_j; 2 \leq j \leq i - 1\} \sim \{x_j; i \leq j \leq n\}
\sim \{y_{n-j}; 0 \leq j \leq n - i - 1\} \sim \{z_j; i + 1 \leq j \leq n\}
\sim \{z_j; 1 \leq j \leq i\} = v.
\]

(15)

Case 5.3: \( i = n \):

\[
P_H(u, v); u = y_1 z_1 \sim \{z_{n-j}; 0 \leq j \leq n - 2\}
\sim \{y_j; 2 \leq j \leq n - 1\} \sim \{x_{n-j}; 0 \leq j \leq n - 1\}
\sim y_n = v.
\]

(16)

Case 6: \( u = y_1 \) and \( v = x_i, \ 1 \leq i \leq n \).

Case 6.1: \( i = 1, 2 \):

\[
P_H(u, v); u = y_1 \sim \{z_j; 1 \leq j \leq n\} \sim \{y_{n-j}; 0 \leq j \leq n - 2\}
\sim \{x_j; i + 1 \leq j \leq n\} \sim \{x_j; 1 \leq j \leq i\} = v.
\]

(17)

Case 6.2: \( 3 \leq i \leq n - 1 \):

\[
P_H(u, v); u = y_1 \sim \{z_j; 1 \leq j \leq n\} \sim \{y_{n-j}; 0 \leq j \leq n - i\}
\sim \{x_j; i + 1 \leq j \leq n\} \sim \{x_j; 1 \leq j \leq i\} = v.
\]
Complexity

Theorem 3. The $n$-dimensional convex polytope $S_n$, with $n \geq 4$, is Hamilton-connected.

Proof. We prove it by definition. This implies that we need to show the existence of Hamiltonian paths between any pair of vertices of $G$.

Let $P_H(u,v)$ be a Hamiltonian path between vertices $u$ and $v$ in $S_n$. By following the labeling of vertices exhibited in Figure 2, we show the existence of Hamiltonian paths between vertices of $S_n$ in a number of cases.
Case 1: \( u = z_i \) and \( v = z_j, \ 2 \leq i \leq n \).

Case 1.1: \( i|2 \):

\[
P_H(u, v) = u \sim z_1 \sim \left\{ w_{j-n} y_j \quad 0 \leq j \leq n - i \right\} \sim \{ y_j; i \leq j \leq n \} \sim \{ z_{n-j}; 0 \leq j \leq n - i \} = v. \tag{30}
\]

Case 1.2: \( i \nmid 2 \):

\[
P_H(u, v) = u \sim z_1 \sim \left\{ y_j; i-1 \leq j \leq n - 1 \right\} \sim \{ x_{n-j}; 1 \leq j \leq n - i + 1 \} \sim \{ w_j; i-1 \leq j \leq n \} \sim x_{i-1} y_i \sim \{ z_{n-j}; 0 \leq j \leq n - i \} = v. \tag{31}
\]

Case 2: \( u = z_1 \) and \( v = y_i, \ 1 \leq i \leq n \).

Case 2.1: \( 1 \leq i \leq n - 1 \):

\[
P_H(u, v) = u \sim z_1 \sim \left\{ y_j; 0 \leq j \leq n - i - 1 \right\} \sim \{ x_j; i+1 \leq j \leq n \} \sim w_1 \sim \{ w_{n-j}; 0 \leq j \leq n - i - 1 \} \sim \{ x_{i-j} w_{i-j}; 0 \leq j \leq i - 2 \} \sim x_1 \sim \{ y_j; 1 \leq j \leq i \} = v. \tag{32}
\]

Case 2.2: \( i = n \):

\[
P_H(u, v) = u \sim z_1 \sim \left\{ w_{n-j}; 0 \leq j \leq n - 2 \right\} \sim \{ x_j; 2 \leq j \leq n - 1 \} \sim \{ y_{n-j}; 1 \leq j \leq n - 2 \} \sim \{ z_j; 2 \leq j \leq n \} \sim y_n = v. \tag{33}
\]

Case 3: \( u = z_1 \) and \( v = x_i, \ 1 \leq i \leq n \).

Case 3.1: \( i = 1 \):

\[
\]
\[
P_H(u, v): u = z_1 \sim \{z_{m-j}: 0 \leq j \leq n - 2\} \sim y_1 \sim \{y_{m-j}: 0 \leq j \leq n - 3\} \sim \{x_j: 3 \leq j \leq n\} \sim w_1 \sim \{w_{m-j}: 0 \leq j \leq n - 2\} \sim x_2 x_1 = v. \\
\]  
(34)

Case 3.2: \(2 \leq i \leq n - 1\):

\[
P_H(u, v): u = z_1 \sim \{z_{m-j}: 0 \leq j \leq n - 2\} \sim \{y_j: 2 \leq j \leq n\} \sim y_1 x_1 \sim \{x_j: 0 \leq j \leq n - i - 1\} \sim \{w_j: i + 1 \leq j \leq n\} \\
\sim w_1 \sim \{w_j, x_j: 2 \leq j \leq i\} = v. \\
\]  
(35)

Case 3.3: \(i = n\):

\[
P_H(u, v): u = \{z_{j}: 1 \leq j \leq n\} \sim \{y_{m-j}: 0 \leq j \leq n - 1\} \sim \{x_j: 1 \leq j \leq n - 1\} \sim \{w_{m-j}: 0 \leq j \leq n - 1\} \sim x_n = v. \\
\]  
(36)

Case 4: \(u = z_1\) and \(v = w_i, 1 \leq i \leq n\).

Case 4.1: \(i = 1\):

\[
P_H(u, v): u = \{z_{j}: 1 \leq j \leq n\} \sim \{y_{m-j}: 0 \leq j \leq n - 1\} \sim \{x_j: 1 \leq j \leq n\} \sim \{w_{m-j}: 0 \leq j \leq n - 1\} = v. \\
\]  
(37)

Case 4.2: \(2 \leq i \leq n - 1\):

\[
P_H(u, v): u = \{z_{j}: 1 \leq j \leq n\} \sim y_n \sim \{y_j: 1 \leq j \leq n - 1\} \sim \{x_j: 1 \leq j \leq n - i - 1\} \sim \{w_j: i + 1 \leq j \leq n\} \\
\sim w_1 x_n \sim \{x_{j-1}, w_j: 2 \leq j \leq i\} = v. \\
\]  
(38)

Case 4.3: \(i = n\):

\[
P_H(u, v): u = \{z_{j}: 1 \leq j \leq n\} \sim \{y_{m-j}: 0 \leq j \leq n - 1\} \sim \{x_j: 1 \leq j \leq n\} \sim \{w_{j}: 1 \leq j \leq n\} = v. \\
\]  
(39)

Case 5: \(u = y_1\) and \(v = z_i, 1 \leq i \leq n\).

Case 5.1: \(i = 1\):

\[
P_H(u, v): u = \{y_{j}: 1 \leq j \leq n - 1\} \sim \{x_{m-j}: 1 \leq j \leq n - 1\} \sim \{w_{j}: 1 \leq j \leq n\} \sim x_n y_n \sim \{z_{m-j}: 0 \leq j \leq n - 1\} = v. \\
\]  
(40)

Case 5.2: \(2 \leq i \leq n - 1\):
\( P_H(u, v): u = \{ y_j: 1 \leq j \leq i \} \sim \{ x_{i-j} w_{i-j}: 0 \leq j \leq i - 1 \} \sim \{ w_{n-j}: 0 \leq j \leq n - i - 1 \} \sim \{ x_j: i + 1 \leq j \leq n \} \)
\( \sim \{ y_{n-j}: 0 \leq j \leq n - i - 1 \} \sim \{ z_j: i + 1 \leq j \leq n \} \)
\( \{ z_j: 1 \leq j \leq i \} = v. \)  \hspace{1cm} (41)

Case 5.3: \( i = n \):

\( P_H(u, v): u = y_1 \sim \{ z_j: 1 \leq j \leq n - 1 \} \sim \{ y_{n-j}: 1 \leq j \leq n - 2 \} \sim \{ x_j: 2 \leq j \leq n - 1 \} \sim \{ w_{n-j}: 0 \leq j \leq n - 1 \} \sim x_1 x_n z_n = v. \)  \hspace{1cm} (42)

Case 6: \( u = y_1 \) and \( v = y_j \), \( 2 \leq i \leq n \).

Case 6.1: \( 2 \leq i \leq n - 1 \):

\( P_H(u, v): u = y_1 \sim \{ z_j: 1 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n - i - 1 \} \sim \{ x_{i-j+1}: 0 \leq j \leq i - 2 \} \sim \{ w_{j}: 3 \leq j \leq i + 1 \} \)
\( \sim \{ w_j x_j: i + 2 \leq j \leq n \} \sim w_1 x_1 w_2 x_2 \sim \{ y_j: 2 \leq j \leq i \} = v. \)  \hspace{1cm} (43)

Case 6.2: \( i = n \):

\( P_H(u, v): u = y_1 z_1 \sim \{ z_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ y_j: 2 \leq j \leq n - 1 \} \sim \{ x_{n-j}: 1 \leq j \leq n - 1 \} \sim \{ w_j: 1 \leq j \leq n \} \sim x_n y_n = v. \)  \hspace{1cm} (44)

Case 7: \( u = y_1 \) and \( v = x_i \), \( 1 \leq i \leq n \).

Case 7.1: \( 1 \leq i \leq n - 1 \):

\( P_H(u, v): u = y_1 z_1 \sim \{ z_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ y_j: 2 \leq j \leq n \} \sim \{ x_{n-j}: 0 \leq j \leq n - i - 1 \} \sim \{ w_j: i + 1 \leq j \leq n \} \)
\( \sim \{ w_j x_j: 1 \leq j \leq i \} = v. \)  \hspace{1cm} (45)

Case 7.2: \( i = n \):

\( P_H(u, v): u = y_1 \sim \{ z_j: 1 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ x_j: 2 \leq j \leq n - 1 \} \sim \{ w_{n-j}: 0 \leq j \leq n - 1 \} \sim x_1 x_n = v. \)  \hspace{1cm} (46)

Case 8: \( u = y_1 \) and \( v = w_i \), \( 1 \leq i \leq n \).

Case 8.1: \( i = 1 \):

\( P_H(u, v): u = y_1 z_1 \sim \{ z_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ y_j: 2 \leq j \leq n \} \sim \{ x_{n-j}: 0 \leq j \leq n - 1 \} \sim \{ w_j: 2 \leq j \leq n \} \sim w_1 = v. \)  \hspace{1cm} (47)

Case 8.2: \( 2 \leq i \leq n \):
\[ P_H(u, v): u = y_1z_1 \sim \{ z_{n-j}; 0 \leq j \leq n - 2 \} - \{ y_j; 2 \leq j \leq n \} - \{ x_{n-j}; 0 \leq j \leq n - i \} - \{ x_{i-j}w_{i-j}; 1 \leq j \leq i - 1 \} \]
\[ \sim \{ w_{n-j}; 0 \leq j \leq n - i \} = v. \]  

Case 9: \( u = x_1 \) and \( v = z_j; 1 \leq i \leq n. \)

9.1: \( i = 1: \)

\[ P_H(u, v): u = x_1 \sim \{ y_j; 1 \leq j \leq n - 1 \} - \{ x_{n-j}; 1 \leq j \leq n - 2 \} - \{ w_j; 2 \leq j \leq n \} - x_1x_ny_n - \{ z_{n-j}; 0 \leq j \leq n - 1 \} = v. \]  

Case 9.2: \( 2 \leq i \leq n - 1: \)

\[ P_H(u, v): u = x_1y_1z_1 \sim \{ z_{n-j}; 0 \leq j \leq n - i - 1 \} - \{ y_j; i + 1 \leq j \leq n \} - \{ x_{n-j}; 0 \leq j \leq n - i - 1 \} - \{ w_j; i + 1 \leq j \leq n \} \]
\[ \sim w_1 - \{ w_jx_j; 2 \leq j \leq i \} - \{ y_{i-j}; 0 \leq j \leq i - 2 \} - \{ z_{j}; 2 \leq j \leq i \} = v. \]  

Case 9.3: \( i = n: \)

\[ P_H(u, v): u = x_1w_1 \sim \{ w_{n-j}; 0 \leq j \leq n - 2 \} - \{ y_j; 2 \leq j \leq n \} - \{ y_{n-j}; 0 \leq j \leq n - 1 \} - \{ z_{n-j}; 1 \leq j \leq n \} = v. \]  

Case 10: \( u = x_1 \) and \( v = y_j; 1 \leq i \leq n. \)

10.1: \( i = 1: \)

\[ P_H(u, v): u = x_1 \sim \{ w_j; 1 \leq j \leq n \} - \{ x_{n-j}; 0 \leq j \leq n - 2 \} - \{ y_j; 2 \leq j \leq n \} - \{ z_{n-j}; 0 \leq j \leq n - 1 \} - y_1 = v. \]  

Case 10.2: \( 2 \leq i \leq n - 1: \)

\[ P_H(u, v): u = x_1y_1z_1 \sim \{ z_{j}; 1 \leq j \leq n \} - \{ y_{n-j}; 0 \leq j \leq n - i - 1 \} - \{ x_{j}w_{i-j}; i + 1 \leq j \leq n \} - \{ w_j; 1 \leq j \leq i \} \]
\[ \sim \{ x_{i-j}; 0 \leq j \leq i - 2 \} - \{ y_j; 2 \leq j \leq i \} = v. \]  

Case 10.3: \( i = n: \)

\[ P_H(u, v): u = x_1y_1z_1 \sim \{ z_{n-j}; 0 \leq j \leq n - 2 \} - \{ y_j; 2 \leq j \leq n - 1 \} - \{ x_{n-j}; 1 \leq j \leq n - 2 \} - \{ w_j; 2 \leq j \leq n \} - w_1x_ny_n = v. \]  

Case 11: \( u = x_1 \) and \( v = x_j; 2 \leq i \leq n. \)

11.1: \( 2 \leq i \leq n - 1: \)
\[ P_H(u, v): u = x_i y_i z_i \sim \{ z_{n-j}: 0 \leq j \leq n-2 \} \sim \{ y_j: 2 \leq j \leq n \} \sim \{ x_{n-j}: 0 \leq j \leq n-i - 1 \} \sim \{ w_j: i+1 \leq j \leq n \} \sim w_i \sim \{ w_j x_j: 2 \leq j \leq i \} = v. \]  

(55)

Case 11.2: \( i = n \):

\[ P_H(u, v): u = x_i y_i \sim \{ z_{n-j}: 0 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ x_j: 2 \leq j \leq n-1 \} \sim \{ w_{n-j}: 0 \leq j \leq n-1 \} \sim x_n = v. \]  

(56)

Case 12: \( u = x_i \) and \( v = w_i, 1 \leq i \leq n \).

Case 12.1: \( i = 1 \):

\[ P_H(u, v): u = x_i y_i \sim \{ z_{n-j}: 0 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n-2 \} \sim \{ x_j: 2 \leq j \leq n-1 \} \sim \{ w_j: 2 \leq j \leq n \} \sim w_i = v. \]  

(57)

Case 12.2: \( 2 \leq i \leq n-1 \):

\[ P_H(u, v): u = x_i y_i \sim \{ z_{n-j}: 0 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n-2 \} \sim \{ x_j: 2 \leq j \leq n-1 \} \sim \{ w_j: i+1 \leq j \leq n \} \sim w_i = v. \]  

(58)

Case 12.3: \( i = n \):

\[ P_H(u, v): u = \{ x_j: 1 \leq j \leq n-1 \} \sim \{ y_{n-j}: 0 \leq j \leq n-1 \} \sim \{ z_j: 1 \leq j \leq n \} \sim y_n x_n \sim \{ w_j: 1 \leq j \leq n \} = v. \]  

(59)

Case 13: \( u = w_i \) and \( v = z_i, 1 \leq i \leq n \).

Case 13.1: \( i = 1 \):

\[ P_H(u, v): u = \{ w_j: 1 \leq j \leq n \} \sim \{ x_{n-j}: 0 \leq j \leq n-1 \} \sim \{ y_j: 1 \leq j \leq n \} \sim \{ z_{n-j}: 0 \leq j \leq n-1 \} = v. \]  

(60)

Case 13.2: \( 2 \leq i \leq n \):

\[ P_H(u, v): u = \{ w_j x_j: 1 \leq j \leq i-1 \} \sim \{ w_j: i \leq j \leq n \} \sim \{ x_{n-j}: 0 \leq j \leq n-i \} \sim \{ y_j: i \leq j \leq n \} \sim \{ y_j: 1 \leq j \leq i-1 \} \sim \{ z_{n-j}: 1 \leq j \leq i-1 \} \sim \{ z_{n-j}: 0 \leq j \leq n-i \} = v. \]  

(61)

Case 14: \( u = w_i \) and \( v = y_i, 1 \leq i \leq n \).

Case 14.1: \( i = 1 \):
\[ P_H(u, v): u = w_1 \sim \{w_{n-j}: 0 \leq j \leq n-2\} \sim \{x_j: 1 \leq j \leq n\} \sim \{y_{n-j}: 0 \leq j \leq n-2\} \sim \{z_j: 2 \leq j \leq n\} \sim z_1 y_1 = v. \] (62)

Case 14.2: \(2 \leq i \leq n-1\):
\[ P_H(u, v): u = \{w_j: 1 \leq j \leq i\} \sim \{w_{i+1}: i+1 \leq j \leq n\} \sim \{x_{n-j}: 0 \leq j \leq n-i-1\} \sim \{y_j: 1 \leq j \leq i-1\} \sim \{z_{n-j}: 0 \leq j \leq n-i\} \sim y_i = v. \] (63)

Case 14.3: \(i = n\):
\[ P_H(u, v): u = \{w_j: 1 \leq j \leq n\} \sim \{x_{n-j}: 0 \leq j \leq n-1\} \sim \{y_j: 1 \leq j \leq n-1\} \sim \{z_{n-j}: 1 \leq j \leq n-1\} \sim y_n y_n = v. \] (64)

Case 15: \(u = w_1\) and \(v = x_i\), \(1 \leq i \leq n\).
Case 15.1: \(i = 1\):
\[ P_H(u, v): u = \{w_j: 1 \leq j \leq n\} \sim \{x_{n-j}: 0 \leq j \leq n-2\} \sim \{y_j: 2 \leq j \leq n\} \sim \{z_{n-j}: 0 \leq j \leq n-1\} \sim y_1 x_1 = v. \] (65)

Existence of Hamiltonian paths between any two vertices of the \(n\)-dimensional convex polytope \(S_n\) completes the proof.

Using Theorems 1 and 3, the following proposition computes the detour index of \(S_n\).

**Proposition 2.** Let \(G = S_n\), where \(n \geq 3\). Then, the detour index of \(G\) is
\[ \omega(G) = \frac{4n(4n-1)^2}{2}. \] (70)

**Proof.** The number of vertices in graph \(G\) is \(4n\). Replacing \(4n\) with \(n\) in Theorem 1 completes the proof. \(\square\)

4. Hamilton Connectivity and the Detour Index of \(K_n\)

This section shows that the \(n\)-dimensional convex polytope \(K_n\) is Hamilton-connected. Its Hamilton connectivity is then used to calculate exact analytical formula of the detour index of \(K_n\).
Malik et al. [25] introduced the family of $n$-dimensional convex polytope $K_n$. We have

$$V(K_n) = \{v_i: w_j; x_j; y_j; z_j: 1 \leq j \leq n\},$$

$$E(K_n) = \{v_jv_{j+1}; w_jw_{j+1}; x_jx_{j+1}; y_jy_{j+1}; z_jz_{j+1}: 1 \leq j \leq n\} \cup \{v_jy_j; v_jw_{j+1}; w_jx_j; x_jy_j; x_jy_{j+1}; y_jz_j: 1 \leq j \leq n\}.$$  

(71)

See Figure 3 to view the $n$-dimensional convex polytope $K_n$.

**Theorem 4.** The graph $n$-dimensional convex polytope $K_n$, with $n \geq 5$, is Hamilton-connected.

**Proof.** We prove this result by definition. For this, we have to show that there exists a Hamiltonian path between any pair of vertices of $K_n$.

Let $P_H(u', v')$ be a Hamiltonian path between vertices $u'$ and $v'$ in $K_n$. Let $V(R_n) = Z \cup Y \cup X \cup V$ such that

$$Z = \{z_1, z_2, \ldots, z_n\}, \quad Y = \{y_1, y_2, \ldots, y_n\}, \quad X = \{x_1, x_2, \ldots, x_n\}, \text{ and } V = \{v_1, v_2, \ldots, v_n\}$.

We divide the construction of Hamiltonian paths in $K_n$ into a number of cases as follows.

Case 1: $u' = z_1$ and $v' = z_i$, $2 \leq i \leq n$.

Case 1.1: $2 \leq i \leq n - 1$:

$$P_H(u', v') = u' = z_1 \sim \{z_{n-j}: 0 \leq j \leq n - i - 1\} \sim \{y_j: i + 1 \leq j \leq n\} \sim \{x_{n-j}: 0 \leq j \leq n - i\} \sim \{w_{n-j}: 0 \leq j \leq i - 3\} \sim \{w_j: 2 \leq j \leq n\} \sim \{v_{n-j}: 0 \leq j \leq n - 1\} \sim w_1x_1y_1z_i \sim \{z_j: 2 \leq j \leq i\} = v'.$$

(72)

Case 1.2: $i = n$:

$$P_H(u', v') = u' = \{z_j: 1 \leq j \leq n - 1\} \sim \{y_{n-j}: 1 \leq j \leq n - 1\} \sim \{x_j: 1 \leq j \leq n - 1\} \sim \{w_{n-j}: 1 \leq j \leq n - 1\} \sim \{v_j: 1 \leq j \leq n\} \sim w_nx_ny_nz_n = v'.$$

(73)

Case 2: $u' = z_1$ and $v' = y_i$, $1 \leq i \leq n$.

Case 2.1: $1 \leq i \leq n - 1$:

$$P_H(u', v') = u' = \{z_j: 1 \leq j \leq n\} \sim \{y_{n-j}: 0 \leq j \leq n - i - 1\} \sim \{x_j: i + 1 \leq j \leq n\} \sim \{w_{n-j}: 0 \leq j \leq n - i - 1\} \sim \{v_j: i + 1 \leq j \leq n\} \sim \{v_j: 1 \leq j \leq i\} \sim \{x_{i-j}: 0 \leq j \leq i - 1\} \sim \{y_j: 1 \leq j \leq i\} = v'.$$

(74)

Case 2.2: $i = n$:

$$P_H(u', v') = u' = z_1 \sim \{z_{n-j}: 0 \leq j \leq n - 2\} \sim \{y_j: 2 \leq j \leq n - 1\} \sim \{x_{n-j}: 1 \leq j \leq n - 2\} \sim \{w_j: 2 \leq j \leq n\} \sim \{v_{n-j}: 0 \leq j \leq n - 1\} \sim w_1x_1y_1z_n = v'.$$

(75)

Case 3: $u' = z_1$ and $v' = x_i$, $1 \leq i \leq n$.

Case 3.1: $1 \leq i \leq n - 1$:
$P_H(u', v')$: $u' = \{z_j: 1 \leq j \leq n\} \sim \{y_{n-j}: 0 \leq j \leq n-1\} \sim \{x_{n-j}: 0 \leq j \leq n-i-1\} \sim \{w_{i-j}: i+1 \leq j \leq n\} \sim v_1$

$\sim \{v_{n-j}: 0 \leq j \leq n-i-1\} \sim \{w_{n-j}v_{n-j}: 0 \leq j \leq i-2\} \sim w_1 \sim \{x_{i-j}: 1 \leq j \leq i\} = v'$.

**Case 3.2: $i = n$:**

$P_H(u', v')$: $u' = \{z_j: 1 \leq j \leq n\} \sim \{y_{n-j}: 0 \leq j \leq n-1\} \sim \{x_{j}: 1 \leq j \leq n-1\} \sim \{w_{n-j}: 1 \leq j \leq n-1\}$

$\sim \{v_{j}: 1 \leq j \leq n\} \sim w_n = v'$.

**Case 4: $u' = z_1$ and $v' = w_i$, $1 \leq i \leq n$.**

**Case 4.1: $1 \leq i \leq n-1$:**

$P_H(u', v')$: $u' = \{z_j: 1 \leq j \leq n\} \sim \{y_{n-j}: 0 \leq j \leq n-1\} \sim \{x_{j}: 1 \leq j \leq n\} \sim \{w_{n-j}: 0 \leq j \leq n-i-1\} \sim \{v_{j}: i+1 \leq j \leq n\}$

$\sim \{v_{j}w_{j}: 1 \leq j \leq i\} = v'$.

**Case 4.2: $i = n$:**

$P_H(u', v')$: $u' = \{z_j: 1 \leq j \leq n\} \sim \{y_{n-j}: 0 \leq j \leq n-1\} \sim \{x_{j}: 1 \leq j \leq n\} \sim \{w_{n-j}: 2 \leq j \leq n-1\} \sim \{v_{j}: 1 \leq j \leq n\}$

$\sim w_{n-1}x_{n-1}w_n = v'$.

**Case 5: $u' = z_1$ and $v' = v_i$, $1 \leq i \leq n$.**

**Case 5.1: $1 \leq i \leq n-1$:**

$P_H(u', v')$: $u' = \{z_j: 1 \leq j \leq n\} \sim \{y_{n-j}: 0 \leq j \leq n-1\} \sim \{x_{j}: 1 \leq j \leq n\} \sim \{w_{n-j}: 0 \leq j \leq n-i\} \sim \{v_{j}: i+1 \leq j \leq n\}$

$\sim \{v_{j}w_{j}: 1 \leq j \leq i-1\} \sim v_i = v'$.
Case 5.2: $i = n$:

\[ P_H(u', v') : u' = \{z_j : 1 \leq j \leq n\} \sim \{y_{n-j} : 0 \leq j \leq n - 1\} \sim \{x_j : 1 \leq j \leq n\} \sim \{w_{n-j} : 0 \leq j \leq n - 1\} \sim \{v_j : 1 \leq j \leq n\} = v'. \] (81)

Case 6: $u' = y_i$ and $v' = z_i$, $1 \leq i \leq n$:

Case 6.1: $1 \leq i \leq n - 1$:

\[ P_H(u', v') : u' = \{y_j x_j : 1 \leq j \leq i\} \sim \{x_j : i + 1 \leq j \leq n - 1\} \sim \{w_{n-j} : 1 \leq j \leq n - 1\} \sim \{v_j : 1 \leq j \leq n\} \sim \{z_j : i + 1 \leq j \leq n\} \sim \{z_j : 1 \leq j \leq i\} = v'. \] (82)

Case 6.2: $i = n$:

\[ P_H(u', v') : u' = y_1 \sim \{z_j : 1 \leq j \leq n - 1\} \sim \{y_{n-j} : 1 \leq j \leq n - 2\} \sim \{x_j : 1 \leq j \leq n - 1\} \sim \{w_{n-j} : 1 \leq j \leq n - 1\} \sim \{v_j : 1 \leq j \leq n\} \sim w_n x_n y_n x_n = v'. \] (83)

Case 7: $u' = y_i$ and $v' = y_i$, $2 \leq i \leq n$:

Case 7.1: $2 \leq i \leq n - 1$:

\[ P_H(u', v') : u' = y_1 \sim \{z_j : 1 \leq j \leq n - 1\} \sim \{y_{n-j} : 0 \leq j \leq n - 2\} \sim \{x_j : i \leq j \leq n\} \sim \{w_{n-j} : 0 \leq j \leq n - 2\} \sim \{v_j : 2 \leq j \leq n\} \sim v_1 w_1 \sim \{x_{i-1} y_j : 2 \leq j \leq i\} = v'. \] (84)

Case 7.2: $i = n$:

\[ P_H(u', v') : u' = y_1 z_1 \sim \{z_{n-j} : 0 \leq j \leq n - 2\} \sim \{y_j : 2 \leq j \leq n - 1\} \sim \{x_{n-j} : 1 \leq j \leq n - 1\} \sim \{w_j : 1 \leq j \leq n - 1\} \sim \{v_{n-j} : 0 \leq j \leq n - 1\} \sim w_n x_n y_n = v'. \] (85)

Case 8: $u' = y_i$ and $v' = x_j$, $1 \leq i \leq n$:

Case 8.1: $1 \leq i \leq n - 1$:

\[ P_H(u', v') : u' = y_1 z_1 \sim \{z_{n-j} : 0 \leq j \leq n - 2\} \sim \{y_j : 2 \leq j \leq n\} \sim \{x_{n-j} : 0 \leq j \leq n - i - 1\} \sim w_{i+1} v_{i+1} \sim \{v_j : i + 2 \leq j \leq n\} \sim \{v_j : 1 \leq j \leq i\} \sim \{w_{n-j} : 0 \leq j \leq i - 1\} \sim \{x_j : 1 \leq j \leq i\} = v'. \] (86)

Case 8.2: $i = n$:
\[ P_H(u', v'): u' = y_1 \sim \{ z_j: 1 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ x_j: 1 \leq j \leq n - 1 \} \sim \{ w_{n-j}: 1 \leq j \leq n - 1 \} \sim \{ v_j: 1 \leq j \leq n \} = v'. \] (87)

Case 9: \( u' = y_1 \) and \( v' = w_i \), \( 1 \leq i \leq n \).

Case 9.1: \( 1 \leq i \leq n - 1 \):

\[ P_H(u', v'): u' = y_1 \sim \{ z_j: 1 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ x_j: 1 \leq j \leq n - 1 \} \sim \{ w_{n-j}: 0 \leq j \leq n - i - 1 \} \sim \{ v_j: i + 1 \leq j \leq n \} \sim \{ v_i w_j: 1 \leq j \leq i \} = v'. \] (88)

Case 9.2: \( i = n \):

\[ P_H(u', v'): u' = y_1 z_1 \sim \{ z_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ y_{n-j}: 2 \leq j \leq n \} \sim \{ x_{n-j}: 0 \leq j \leq n - 1 \} \sim \{ w_{n-j}: 1 \leq j \leq n - 1 \} \sim \{ v_{n-j}: 0 \leq j \leq n - 1 \} \sim \{ v_i: 0 \leq j \leq n - 1 \} = v'. \] (89)

Case 10: \( u' = y_1 \) and \( v' = v_i \), \( 1 \leq i \leq n \).

Case 10.1: \( 1 \leq i \leq n - 1 \):

\[ P_H(u', v'): u' = y_1 \sim \{ z_j: 1 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ x_j: 1 \leq j \leq n - 1 \} \sim \{ w_{n-j}: 0 \leq j \leq n - i \} \sim \{ v_j: i + 1 \leq j \leq n \} \sim \{ v_i w_j: 1 \leq j \leq i - 1 \} \sim v = v'. \] (90)

Case 10.2: \( i = n \):

\[ P_H(u', v'): u' = y_1 \sim \{ z_j: 1 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ x_j: 1 \leq j \leq n \} \sim \{ w_{n-j}: 0 \leq j \leq n - 1 \} \sim \{ v_j: 1 \leq j \leq n \} = v'. \] (91)

Case 11: \( u' = x_1 \) and \( v' = z_i \), \( 1 \leq i \leq n \).

Case 11.1: \( 1 \leq i \leq n - 1 \):

\[ P_H(u', v'): u' = x_1 \sim \{ y_j x_j: 2 \leq j \leq i \} \sim \{ w_{i-j}: 0 \leq j \leq i - 2 \} \sim \{ w_{n-j}: 0 \leq j \leq n - i - 2 \} \sim \{ v_j: i + 2 \leq j \leq n \} \sim \{ v_j: 1 \leq j \leq i + 1 \} \sim w_{i+1} \sim \{ x_j: i + 1 \leq j \leq n \} \sim y_1 \sim \{ y_{n-j}: 0 \leq j \leq n - i - 1 \} \sim \{ z_j: i + 1 \leq j \leq n \} \sim \{ z_j: 1 \leq j \leq i \} = v'. \] (92)

Case 11.2: \( i = n \):
\[ P_H(u', v') : \begin{align*}
u' &= x_1 w_1 \sim \{ v_j : 1 \leq j \leq n \} \sim \{ w_{n-j} : 0 \leq j \leq n-2 \} \sim \{ x_j : 2 \leq j \leq n \} \sim \{ y_{n-j} : 0 \leq j \leq n-1 \} \\
&\quad \sim \{ z_j : 1 \leq j \leq n \} = v'. \tag{93} \end{align*} \]

Case 12: \( u' = x_1 \) and \( v' = y_i, 1 \leq i \leq n \).

Case 12.1: \( i = 1 \):

\[ P_H(u', v') : u' = x_1 w_1 v_1 \sim \{ v_{n-j} : 0 \leq j \leq n-2 \} \sim \{ w_j : 2 \leq j \leq n \} \sim \{ x_{n-j} : 0 \leq j \leq n-2 \} \sim \{ y_j : 2 \leq j \leq n \} \\
&\quad \sim \{ z_j : 0 \leq j \leq n-1 \} = y_1 = v'. \tag{94} \]

Case 12.2: \( 2 \leq i \leq n-1 \):

\[ P_H(u', v') : u' = x_1 w_1 v_1 \sim \{ v_{n-j} : 0 \leq j \leq n-2 \} \sim \{ w_j : 2 \leq j \leq n \} \sim \{ x_{n-j} : 0 \leq j \leq n-i \} \sim \{ y_j : i+1 \leq j \leq n \} \\
&\quad \sim \{ z_j : 0 \leq j \leq n-1 \} = y_1 = x_1 \sim \{ y_{j-1} : 2 \leq j \leq i-1 \} \sim y_i = v'. \tag{95} \]

Case 12.3: \( i = n \):

\[ P_H(u', v') : u' = x_1 w_1 v_1 \sim \{ z_{n-j} : 0 \leq j \leq n-2 \} \sim \{ y_j : 2 \leq j \leq n-1 \} \sim \{ x_{n-j} : 1 \leq j \leq n-2 \} \sim \{ w_j : 2 \leq j \leq n-1 \} \\
&\quad \sim \{ v_{n-j} : 0 \leq j \leq n-2 \} \sim w_1 v_1 w_n x_n y_n = v'. \tag{96} \]

Case 13: \( u' = x_1 \) and \( v' = x_i, 2 \leq i \leq n \).

Case 13.1: \( i = 2 \):

\[ P_H(u', v') : u' = x_1 \sim \{ y_j : 1 \leq j \leq n-1 \} \sim \{ z_{n-j} : 1 \leq j \leq n-1 \} \sim z_n y_n \sim \{ x_{n-j} : 0 \leq j \leq n-3 \} \sim w_3 v_3 \sim \{ v_j w_j : 4 \leq j \leq n \} \\
&\quad \sim v_1 w_1 v_2 w_2 x_2 = v'. \tag{97} \]

Case 13.2: \( 3 \leq i \leq n-1 \):

\[ P_H(u', v') : u' = x_1 y_1 \sim \{ y_j x_j : 2 \leq j \leq i-1 \} \sim \{ y_j : i \leq j \leq n-1 \} \sim \{ z_{n-j} : 1 \leq j \leq n-1 \} \sim z_n y_n \\
&\quad \sim \{ x_{n-j} : 0 \leq j \leq n-i-1 \} \sim w_{i+1} v_{i+1} \sim \{ v_j w_j : i+2 \leq j \leq n \} \sim \{ v_j w_j : 1 \leq j \leq i \} \sim x_i = v'. \tag{98} \]

Case 13.3: \( i = n \):

\[ P_H(u', v') : u' = x_1 w_1 \sim \{ v_j : 1 \leq j \leq n \} \sim \{ w_{n-j} : 0 \leq j \leq n-2 \} \sim \{ x_j : 2 \leq j \leq n-1 \} \sim \{ y_{n-j} : 1 \leq j \leq n-1 \} \\
&\quad \sim \{ z_j : 1 \leq j \leq n \} \sim y_n x_n = v'. \tag{99} \]

Case 14: \( u' = x_1 \) and \( v' = w_i, 1 \leq i \leq n \).
\[ P_H(u', v'): u' = x_1y_1 \sim \{ z_j: 1 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ x_j: 2 \leq j \leq n \} \sim \{ w_{n-j}: 0 \leq j \leq n - i - 1 \} \sim \{ v_j: i + 1 \leq j \leq n \} \sim \{ v, w_j: 1 \leq j \leq i \} = v'. \] 

(100)

Case 14.2: \( i = n \):

\[ P_H(u', v'): u' = x_1w_1 \sim \{ v_j: 1 \leq j \leq n \} \sim \{ w_{n-j}: 1 \leq j \leq n - 2 \} \sim \{ x_j: 2 \leq j \leq n - 1 \} \sim \{ y_{n-j}: 1 \leq j \leq n - 1 \} \sim \{ z_j: 1 \leq j \leq n \} \sim y_n x_n w_n = v'. \] 

(101)

Case 15: \( u' = x_1 \) and \( v' = v_i, 1 \leq i \leq n \).

Case 15.1: \( 1 \leq i \leq n - 1 \):

\[ P_H(u', v'): u' = x_1y_1 \sim \{ z_j: 1 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ x_j: 2 \leq j \leq n \} \sim \{ w_{n-j}: 0 \leq j \leq n - i \} \sim \{ v_j: i + 1 \leq j \leq n \} \sim \{ v, w_j: 1 \leq j \leq i - 1 \} \sim v_i = v'. \] 

(102)

Case 15.2: \( i = n \):

\[ P_H(u', v'): u' = x_1y_1 \sim \{ z_j: 1 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ x_j: 2 \leq j \leq n \} \sim \{ w_{n-j}: 0 \leq j \leq n - 1 \} \sim \{ v_j: 1 \leq j \leq n \} = v'. \] 

(103)

Case 16: \( u' = w_1 \) and \( v' = z_i, 1 \leq i \leq n \).

Case 16.1: \( 1 \leq i \leq n - 1 \):

\[ P_H(u', v'): u' = w_1v_1 \sim \{ y_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ w_j: 2 \leq j \leq n \} \sim x_n \sim \{ x_j: 1 \leq j \leq i \} \sim \{ y_{i-j}: 0 \leq j \leq i - 1 \} \sim \{ y_{n-j}: 0 \leq j \leq n - i - 1 \} \sim \{ z_j: i + 1 \leq j \leq n \} \sim \{ z_j: 1 \leq j \leq i \} = v'. \] 

(104)

Case 16.2: \( i = n \):

\[ P_H(u', v'): u' = w_1 \sim \{ v_j: 1 \leq j \leq n \} \sim \{ w_{n-j}: 0 \leq j \leq n - 2 \} \sim \{ x_j: 2 \leq j \leq n \} \sim \{ y_{n-j}: 0 \leq j \leq n - 2 \} \sim x_1y_1 \sim \{ z_j: 1 \leq j \leq n \} = v'. \] 

(105)

Case 17: \( u' = w_1 \) and \( v' = y_j, 1 \leq i \leq n \).

Case 17.1: \( 1 \leq i \leq n - 1 \):
\begin{align}
P_H(u', v') : u' = w_1 w_n v_1 & \sim \{v_{n-j} : 0 \leq j \leq n-2\} \sim \{w_j : 2 \leq j \leq n\} \sim \{x_{n-j} y_{n-j} : 1 \leq j \leq n-i-1\} \\
& \sim \{x_{n-i} : 0 \leq i \leq 1\} \sim x_n y_n \sim \{z_{n-j} : 0 \leq j \leq n-1\} \sim \{y_j : 1 \leq j \leq n-1\} = v'. \quad (106)
\end{align}

Case 17.2: \(i = n\):

\begin{align}
P_H(u', v') : u' = w_1 v_1 & \sim \{v_{n-j} : 0 \leq j \leq n-2\} \sim \{w_j : 2 \leq j \leq n\} \sim \{x_{n-j} : 0 \leq j \leq n-1\} \sim \{y_j : 1 \leq j \leq n-1\} \\
& \sim \{z_{n-j} : 1 \leq j \leq n-1\} \sim z_n y_n = v'. \quad (107)
\end{align}

Case 18: \(u' = w_1\) and \(v' = x_i, 1 \leq i \leq n\).

Case 18.1: \(1 \leq i \leq n-1\):

\begin{align}
P_H(u', v') : u' = w_1 v_1 & \sim \{v_{n-j} : 0 \leq j \leq n-2\} \sim \{w_j : 2 \leq j \leq n\} \sim \{x_{n-j} : 0 \leq j \leq n-i-1\} \sim \{y_j : i+1 \leq j \leq n\} \\
& \sim \{z_{n-j} : 0 \leq j \leq n-1\} \sim \{y_j x_j : 1 \leq j \leq i\} = v'. \quad (108)
\end{align}

Case 18.2: \(i = n\):

\begin{align}
P_H(u', v') : u' = w_1 & \sim \{v_j : 1 \leq j \leq n\} \sim \{w_{n-j} : 0 \leq j \leq n-2\} \sim \{x_j : 2 \leq j \leq n-1\} \sim \{y_{n-j} : 1 \leq j \leq n-2\} \\
& \sim x_1 y_1 \sim \{z_j : 1 \leq j \leq n\} \sim y_n x_n = v'. \quad (109)
\end{align}

Case 19: \(u' = w_1\) and \(v' = w_i, 2 \leq i \leq n\).

Case 19.1: \(2 \leq i \leq n-1\):

\begin{align}
P_H(u', v') : u' = w_1 x_1 y_1 & \sim \{z_j : 1 \leq j \leq n\} \sim \{y_{n-j} : 0 \leq j \leq n-2\} \sim \{x_j : 2 \leq j \leq n\} \sim \{w_{n-j} : 0 \leq j \leq n-i-1\} \\
& \sim \{v_j : i+1 \leq j \leq n\} \sim v_1 \sim \{v_j w_j : 2 \leq j \leq n\} = v'. \quad (110)
\end{align}

Case 19.2: \(i = n\):

\begin{align}
P_H(u', v') : u' = w_1 x_1 y_1 z_1 & \sim \{z_{n-j} : 0 \leq j \leq n-2\} \sim \{y_j : 2 \leq j \leq n\} \sim \{x_{n-j} : 0 \leq j \leq n-2\} \\
& \sim \{w_j : 2 \leq j \leq n-1\} \sim \{v_{n-j} : 1 \leq j \leq n-1\} \sim v_n w_n = v'. \quad (111)
\end{align}

Case 20: \(u' = w_1\) and \(v' = v_i, 1 \leq i \leq n\).

Case 20.1: \(i = 1\):
\[P_H(u', v') : u' = w_i x_i y_1 \sim \{z_j : 1 \leq j \leq n\} \sim \{y_{n-j} : 0 \leq j \leq n-2\} \sim \{x_j : 2 \leq j \leq n\} \sim \{w_{n-j} : 0 \leq j \leq n-2\} \sim \{v_j : 2 \leq j \leq n\} \sim v_1 = v' \]  

(112)

Case 20.2: \(2 \leq i \leq n-1:\)

\[P_H(u', v') : u' = w_i x_i y_1 \sim \{z_j : 1 \leq j \leq n\} \sim \{y_{n-j} : 0 \leq j \leq n-2\} \sim \{x_j : 2 \leq j \leq n\} \sim \{w_{n-j} : 0 \leq j \leq n-i\} \sim \{v_j : i+1 \leq j \leq n\} \sim v_1 \sim \{v_j w_j : 2 \leq j \leq i-1\} \sim v_i = v'. \]  

(113)

Case 20.3: \(i = n:\)

\[P_H(u', v') : u' = w_i x_i y_1 \sim \{z_{n-j} : 0 \leq j \leq n-2\} \sim \{y_j : 2 \leq j \leq n\} \sim \{x_{n-j} : 0 \leq j \leq n-2\} \sim \{w_j : 2 \leq j \leq n\} \sim \{v_j : 1 \leq j \leq n\} = v'. \]  

(114)

Case 21: \(u' = v_1\) and \(v' = z_i, 1 \leq i \leq n.\)

Case 21.1: \(1 \leq i \leq n-1:\)

\[P_H(u', v') : u' = v_1 \sim \{v_{n-j} : 0 \leq j \leq n-2\} \sim \{w_j : 1 \leq j \leq n\} \sim \{x_{n-j} : 0 \leq j \leq n-i-1\} \sim \{y_j : i+2 \leq j \leq n\} \sim \{y_j x_j : 1 \leq j \leq i\} \sim y_{i+1} \sim \{z_j : i+1 \leq j \leq n\} \sim \{z_j : 1 \leq j \leq i\} = v'. \]  

(115)

Case 21.2: \(i = n:\)

\[P_H(u', v') : u' = \{v_j : 1 \leq j \leq n\} \sim \{w_{n-j} : 0 \leq j \leq n-1\} \sim \{x_j : 1 \leq j \leq n\} \sim \{y_{n-j} : 0 \leq j \leq n-1\} \sim \{z_j : 1 \leq j \leq n\} = v'. \]  

(116)

Case 22: \(u' = v_1\) and \(v' = y_i, 1 \leq i \leq n.\)

Case 22.1: \(1 \leq i \leq n-1:\)

\[P_H(u', v') : u' = v_1 \sim \{v_{n-j} : 0 \leq j \leq n-2\} \sim \{w_j : 1 \leq j \leq n\} \sim \{x_{n-j} : 0 \leq j \leq n-i\} \sim \{y_j : i+1 \leq j \leq n\} \sim \{z_{n-j} : 0 \leq j \leq n-1\} \sim \{y_j x_j : 1 \leq j \leq i-1\} \sim y_i = v'. \]  

(117)

Case 22.2: \(i = n:\)

\[P_H(u', v') : u' = v_1 \sim \{v_{n-j} : 0 \leq j \leq n-2\} \sim \{w_j : 1 \leq j \leq n\} \sim \{x_{n-j} : 0 \leq j \leq n-1\} \sim \{y_j : 1 \leq j \leq n-1\} \sim \{z_{n-j} : 1 \leq j \leq n\} \sim z_{n'n} = v'. \]  

(118)

Case 23: \(u' = v_1\) and \(v' = x_i, 1 \leq i \leq n.\)

Case 23.1: \(1 \leq i \leq n-1:\)
\[ P_H(u', v'): u' = v_1 \sim \{v_{n-j}: 0 \leq j \leq n-2\} \sim \{w_j: 1 \leq j \leq n\} \]
\[ \sim \{x_{n-j}: 0 \leq j \leq n-i-1\} \sim \{y_j: i+1 \leq j \leq n\} \sim \{z_{n-j}: 0 \leq j \leq n-1\} \sim \{y_j x_j: 1 \leq j \leq i\} = v'. \] (119)

Case 23.2: \( i = n \):

\[ P_H(u', v'): u' = \{v_j: 1 \leq j \leq n\} \sim \{w_{n-j}: 0 \leq j \leq n-1\} \sim \{x_j: 1 \leq j \leq n-1\} \sim \{y_{n-j}: 1 \leq j \leq n-1\} \]
\[ \sim \{z_j: 1 \leq j \leq n\} \sim y_n x_n = v'. \] (120)

Case 24: \( u' = v_1 \) and \( v' = w_i \), \( 1 \leq i \leq n \).

Case 24.1: \( 1 \leq i \leq n-1 \):

\[ P_H(u', v'): u' = \{v_j w_j: 1 \leq j \leq i-1\} \sim \{v_j: i \leq j \leq n\} \sim \{w_{n-j}: 0 \leq j \leq n-i-1\} \sim \{x_j: i+1 \leq j \leq n\} \]
\[ \sim \{y_{n-j}: 0 \leq j \leq n-2\} \sim \{z_j: 2 \leq j \leq n\} \sim z_1 y_1 \sim \{x_j: 1 \leq j \leq i\} \sim w_1 = v'. \] (121)

Case 24.2: \( i = n \):

\[ P_H(u', v'): u' = v_1 \sim \{v_{n-j}: 0 \leq j \leq n-2\} \sim \{w_j: 1 \leq j \leq n-1\} \sim \{x_{n-j}: 1 \leq j \leq n-1\} \sim \{y_j: 1 \leq j \leq n-1\} \]
\[ \sim \{z_{n-j}: 1 \leq j \leq n-1\} \sim z_n y_n x_n w_n = v'. \] (122)

Case 25: \( u' = v_1 \) and \( v' = v_j \), \( 2 \leq i \leq n \).

Case 25.1: \( 2 \leq i \leq n-1 \):

\[ P_H(u', v'): u' = v_1 \sim \{v_{n-j}: 0 \leq j \leq n-i-1\} \sim \{w_j: i \leq j \leq n\} \sim \{x_{n-j}: 0 \leq j \leq n-2\} \sim \{y_j: 2 \leq j \leq n\} \]
\[ \sim \{z_{n-j}: 0 \leq j \leq n-1\} \sim y_1 x_1 \sim \{w_{n-j}: 2 \leq j \leq i\} = v'. \] (123)

Case 25.2: \( i = n \):

\[ P_H(u', v'): u' = v_1 w_1 x_1 y_1 \sim \{z_j: 1 \leq j \leq n\} \sim \{y_{n-j}: 0 \leq j \leq n-2\} \sim \{x_j: 2 \leq j \leq n\} \sim \{w_{n-j}: 0 \leq j \leq n-2\} \]
\[ \sim \{v_j: 2 \leq j \leq n\} = v'. \] (124)

Existence of Hamiltonian paths between any two vertices of the \( n \)-dimensional convex polytope \( K_n \) completes the proof. \( \square \)

Using Theorems 1 and 4, the following proposition computes the detour index of \( K_n \).

**Proposition 3.** Let \( G = K_n \), where \( n \geq 3 \). Then the detour index of \( G \) is

\[ \omega(G) = \frac{5n(5n-1)^2}{2}. \] (125)
Proof. The number of vertices in graph $G$ is $5n$. Replacing $5n$ with $n$ in Theorem 1 completes the proof.

5. A Family of Non-Hamilton-connected Convex Polytopes

For $n \geq 4$, by $B_n$, we denote the graph of convex polytope defined in [24] which consists of $n$ number of 3-sided faces, $2n$ number of 4-sided faces, and $n$ number of pentagonal faces, see Figure 4.

Mathematically, the vertex set of $B_n$ consists of four layers of vertices, i.e., $v_i$, $w_i$, $x_i$, $y_i$, and $z_i$. That is to say that $V(B_n) = \{v_i, w_i, x_i, y_i, z_i : 0 \leq i \leq n-1\}$. Accordingly, the edge set of $B_n$ is as follows:

$$E(B_n) = \{v_iv_{i+1}, v_iw_i, w_ix_0, w_ix_i, x_0x_i, x_iy_0, y_0y_i, y_iy_{i+1}, y_0z_i, z_iz_{i+1} : 0 \leq i \leq n-1\}.$$ (126)

Note that arithmetics in the subscripts is performed modulo $n-1$. Layer comprising vertices $v_i (0 \leq i \leq n-1)$ is called the first layer. Similarly, layers composed by vertices $w_i (0 \leq i \leq n-1)$, $x_i (0 \leq i \leq n-1)$, $y_i (0 \leq i \leq n-1)$, and $z_i (0 \leq i \leq n-1)$ are called second layer, third layer, fourth layer, and fifth layer, respectively. Note that first, second, fourth, and fifth layers form cycles of length $n$.

The following result shows that the infinite family $B_n$ of convex polytopes is non-Hamilton-connected.

**Theorem 5.** The $n$-dimensional convex polytope $B_n$ with $n \geq 4$, is non-Hamilton-connected.

**Proof.** It is enough to show that there exist two vertices in the $n$-dimensional convex polytope $B_n$ such that no Hamiltonian path exists between them.

It is easy to see that there exists no Hamiltonian path between any two vertices at distance two on the outer layer, i.e., between $z_i$ and $z_{i+2}$ if $n$ is even and between consecutive vertices on the outer layer, i.e., $z_i$ and $z_{i+1}$ if $n$ is odd, where $0 \leq i \leq n-1$ such that subscripts are taken modulo $n-1$. This shows that the $n$-dimensional convex polytope $B_n$ is non-Hamilton-connected. \qed

6. Conclusion

Determining whether or not a graph is Hamilton-connected is NP-complete. Thus, it is natural to study the Hamilton-connectivity of infinite families of graphs. In this paper, we have studied Hamilton-connectivity of certain infinite families of convex polytopes. We construct three infinite families of convex polytopes which have been shown to be Hamilton-connected. Moreover, one infinite family is shown to be non-Hamilton-connected which shows that not all the convex polytopes are Hamilton-connected. As a by-product, we compute the detour index of Hamilton-connected families of convex polytopes.

As future directions, we propose the following problems:

**Problem 1.** Is there any other way to show Hamilton-connectivity of a given graph?

**Problem 2.** Baca [18] introduced a family $D_n$ of convex polytopes. Determine whether $D_n$ is Hamilton-connected.

**Problem 3.** Imran et al. [57] introduced the family $T_n$ of convex polytopes. Determine whether or not $T_n$ is Hamilton-connected.

Data Availability

No data were used to support the findings of the study.

Conflicts of Interest

The authors declare that they have no conflicts of interest regarding the publication of the article.

Authors’ Contributions

S. H. and A. K. devised the methodology and acquired funding. S. H. and S. K. carried out the formal analysis and data curation. S. H. and J.-B. L. wrote the original draft, reviewed the writing, and edited the manuscript. S. H., A. K.,
and J.-B. L. proofread the manuscript before its final submission. S. H. and A. K. contributed equally to this work.

Acknowledgments

Sakander Hayat was grateful to Dr. Muhammad Imran for providing the registration details to Mayura software. Sakander Hayat was supported by the Higher Education Commission, Pakistan, under Grant no. 20-11682/NRPU/R&D/HEC/2020. J.-B. Liu was supported by Anhui Provincial Natural Science Foundation under Grant 200808501 and Natural Science Fund of Education Department of Anhui province under Grant no. KJ2020A0478.

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