Lambada, Haskell as a better Java

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Abstract
The Lambada framework provides facilities for fluid interoperation between Haskell (currently both Hugs and GHC using non-Haskell98 extensions) and Java. Using Lambada, we can call Java methods from Haskell, and have Java methods invoke Haskell functions. The framework rests on the Java Native Interface (JNI). The Lambada release includes a tool for generating IDL from Java .class files (using reflection), which is fed into our existing HDirect to generate Haskell-callable stubs.

1 Introduction

It goes without saying that the ability to interact with other languages is of vital importance for the long-term survival of any programming language, in particular for niche languages such as Haskell or ML. Java is an interesting partner for interoperation for a number of reasons:

• Java comes with an large set of stable, and usually well-documented and well-designed libraries and APIs.
• The interaction with Java via the Java Native Interface (JNI) [8,4] effectively allows us to script the underlying JVM. This makes bidirectional interoperability with Java very flexible and vendor independent.
• Because of Java’s platform independence, we hope that Lambada gains wider acceptance than our previous work on COM.

Notwithstanding all these advantages, using JNI in its most primitive form is tedious and error prone. Using raw JNI is like assembly programming on the JVM. Our mission is to make interoperation between Java and Haskell via JNI as convenient as programming in Java directly.

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The architecture of Lambada borrows heavily from our previous work on interfacing Haskell and COM [6,2,3], in particular the binding for Automation [7]. We refer the reader to those papers for a discussion on the rationale behind the object encoding used for Lambada.

2 The Java Native Interface

Simply said, JNI provides a number of invocation functions to initialize and instantiate a Java Virtual Machine (JVM), and defines a COM interface to the JVM (the JavaVM interface), and to individual threads that run inside a particular JVM (the JNIEnv interface). The JNIEnv interface contains roughly 230 methods to interact with objects that live in its associated thread.

2.1 Calling C From Java and vice versa

To explain the way C and Java interact via JNI we use a slightly perverse version of Hello World! where a Java method calls a C procedure that calls back into Java to print Hello World!.

To use foreign functions in Java we declare a method as native. Class Hello imports a native method cayHello from the HelloImpl DLL [1] and defines a normal static Java method sayHello that will print the greeting on the standard output:

```java
class Hello {
    native cayHello ();
    static void sayHello () {
        System.out.println("Hello World");
    }
    static {
        System.loadLibrary ("HelloImpl");
    }
}
```

In Java we cannot just call entries of an arbitrary DLL; Java dictates how native methods should look like. For this example HelloImpl.dll should have entry Java_Hello_cayHello with signature [2]:

```c
JNIEXPORT void JNICALL Java_Hello_cayHello ([in]JNIEnv* env, [in]jobject this)
```

---

1 A dynamic-link library (DLL) is an executable file that acts as a shared library of functions on the Win32 platform.

2 Throughout this paper we specify the C signatures for the different JNI procedures using IDL attributes such as [in] and [out].
In general, each native method implementation gets an additional JNIEnv argument, and a jobject object reference (or a jclass class object for static methods). The JNIEnv argument is an interface pointer through which the native method can communicate with Java, for instance to access fields of the object, or to call back into the JVM. The jobject (or jclass) reference corresponds to the this pointer (or class object) in Java.

Procedure Java_Hello_cayHello first gets the class object for the Hello class, then gets the methodID for the static void sayHello() method, and finally calls it:

```c
jclass cls = (*env) -> GetObjectClass (env, this);
jmethodID mid = (*env) -> GetStaticMethodID (env, cls, "sayHello" "()V");
(*env) -> callStaticVoidMethod (env, cls, mid);
```

In the above example the initiative came from the Java side. It is also possible to start on the C side.

The JNI_CreateJavaVM function takes an initialization structure and loads and initializes a JVM instance. It returns both a JavaVM interface pointer to the JVM and a JNIEnv interface pointer to the main thread running in that JVM instance:

```c
jint JNI_CreateJavaVM ([out]JavaVM** v, [out]JNIEnv** e, [in]JavaVMInitArgs* a);
```

To call sayHello in this manner, we first load and initialize a JVM using JNI_CreateJavaVM to obtain a JNIEnv pointer, then call the sayHello method exactly as we did before except that we get the class object via the FindClass function and finish by destroying the JVM:

```c
JNIEnv* env;
JavaVM* jvm;
JavaVMInitArgs args;

new_InitArgs (&args);
JNI_CreateJavaVM (&jvm, &env, &args);
jclass cls = (*env) -> FindClass (env, "Hello");
jmethodID mid = (*env) -> GetStaticMethodID (env, cls, "sayHello" "(V");
(*env) -> callStaticVoidMethod (env, cls, mid);
(*jvm) -> DestroyJavaVM(jvm);
```

Even though we have left out many details (including dealing with errors) it is clear that interacting with Java at this level of abstraction is quite mind numbing. In particular, the fact that we have to thread around the JNIEnv pointer in every call is something we would like to get rid of. In addition,
encoding of the result-type in method calls (such as callStaticVoidMethod) and the distinction between static and instance methods is something we want to hide to the user.

3 Essential FFI in Four Easy Steps

We will introduce some basic JNI concepts that underpin our JNI binding for Haskell while reviewing the basics of the Hugs/GHC foreign function interface. For a more thorough explanation of the FFI we refer to our ICFP99 paper [2].

3.1 Importing Function Pointers

In order to call a foreign function from Haskell, we import externally defined functions into Haskell, either by static linking or by dynamic linking.

For example, we can statically import the LoadLibraryA entry of the Microsoft Windows kernel32 DLL

HMODULE LoadLibraryA([in,string]char*);

by giving the following foreign import declaration:

type CString = Addr

type HMODULE = Int

foreign import stdcall "kernel32"
  "LoadLibraryA" primLoadLibrary :: CString -> IO HMODULE

This foreign declaration defines a Haskell function primLoadLibrary, that takes the address of a null terminated string which contains the filename of an executable module, dynamically loads the module into memory and returns a handle to the requested module.

As its name suggests, function primLoadLibrary is rather primitive. Using the functions marshallString and freeCString we can easily write a friendlier version that takes a normal Haskell string as its argument:

marshallString :: String -> CString
freeCString :: CString -> IO ()

loadLibrary :: FilePath -> IO HMODULE
loadLibrary = \
  do{ s <- marshallString p;
       h <- primLoadLibrary s;
       freeCString s;
       return h;
The kernel32 DLL also exports a function `GetProcAddress` that given a module handle and a pointer to a null-terminated string containing the desired function name returns the address of that function.

```
FARPROC GetProcAddressA ([in]HMODULE, [in,string]char*);
```

To use function `GetProcAddress` from Haskell we declare the following foreign import:

```
type FARPROC = Addr

foreign import "kernel32" "GetProcAddressA"
    primGetProcAddress :: HMODULE -> CString -> IO FARPROC
```

Again we can easily construct a Haskell-friendly version of `primGetProcAddress` that takes a normal Haskell string as its argument:

```
getProcAddress :: HMODULE -> String -> IO FARPROC
getProcAddress = \h -> \n -> do{
    pn <- marshallString n;
    pf <- primGetProcAddress h pn;
    freeCString pn;
    return pf;
}
```

The explicit freeing of `CString`s is already getting tedious but don’t despair; in section 3.2 we will instruct the Haskell garbage collector to do this automatically.

When interfacing with Java using JNI, we have to deal with different JVM implementations and different versions of the same JVM. Therefore, we don’t want to commit ourselves to a particular VM or version too early. The `loadLibrary` and `GetProcAddress` functions together with *dynamic* foreign function import allow us to delay the choice of a particular JVM to the very last moment.

The dynamic variation of the foreign import declaration defines a function that may be used to unmarshall an external function pointer into a Haskell function.

We can use a function pointer that represents the `JNI_CreateJavaVM` DLL entry by declaring the dynamic foreign import `mkPrimCreateJavaVM`. This function will coerce the low-level function pointer into the corresponding Haskell function:

```
type JavaVM = Addr

type JNIEnv = Addr
```
To write a flexible version of `createJavaVM` we need several helper functions. Function `newJavaVMInitArgs` creates a default JVM initialization structure; it comes with a corresponding free routine `freeJavaVMInitArgs`:

```haskell
newJavaVMInitArgs :: IO Addr
freeJavaVMInitArgs :: Addr -> IO ()
```

Function `newResultAddr` allocates space for a return value; it also comes with a corresponding free routine `freeResultAddr`:

```haskell
newResultAddr :: IO Addr
freeResultAddr :: Addr -> IO ()
```

And finally, a function `deref` to shorten a pointer indirection:

```haskell
deref :: Addr -> IO Addr
```

Using these helper functions, we can complete the tedious code for `createJavaVM`:

```haskell
type JavaVM = Addr
type JNIEnv = Addr

createJavaVM :: FilePath
             -> IO (JavaVM, JNIEnv)
createJavaVM = \p -> do{  
    h <- loadLibrary p; 
    a <- getprocAddress h "JNI_CreateJavaVM"; 
    pArgs <- newJavaVMInitArgs; 
    pJavaVM <- newResultAddr; 
    pJNIEnv <- newResultAddr; 
    mkPrimCreateJavaVM a pJavaVM pJNIEnv pArgs; 
    freeJavaVMInitArgs pArgs; 
    jvm <- deref pJavaVM; 
    freeResultAddr pJavaVM; 
    jnienv <- deref pJNIEnv; 
    freeResultAddr pJNIEnv; 
    return (jvm,jnienv);  
}
```

There are numerous other situations where we want to call dynamic function pointers. For instance both the `JavaVM` and the `JNIEnv` pointers are a double indirection to a table of function pointers (a vtable). The function `deref` we used earlier and the function `getProcAt t i` that returns the i-the function
pointer in table t make it quite easy to call vtable entries. We will give an example of this in section B.3.

3.2 Interfacing With the GC Using Foreign Objects

One of the nice things about languages such as Haskell and Java is that the garbage collector silently takes care of freeing unused resources. It seems like a great loss therefore that when we import foreign functions, we import the hassle of explicit resource management as well.

What we would like to do instead is to have a hook into the Haskell garbage collector so that we can tell it how to free external resources once they become garbage. As usual, an extra level of indirection does the job. A foreign object is a smart pointer to the outside world:

\[
\text{data ForeignObject}\\
\text{new_ForeignObject :: Addr} \to (\text{Addr} \to \text{IO} () ) \to \text{IO ForeignObject}\\
\]

For example, we can define a smart String marshaller that frees the string automatically once it has become garbage (to avoid confusion, the ‘dumb’ version of marshallString is qualified with prefix Prim):

\[
\text{type CString = ForeignObject}\\
\text{marshallString :: String} \to \text{IO CString}\\
\text{marshallString = \s} \to \text{do}\{\\
\text{ps <- Prim.marshallString s;}\\
\text{newForeignObject Prim.freeString ps; }\\
\}\n\]

Note that imported functions can directly take a ForeignObject as argument. Extracting the underlying Addr first, would be rather dangerous as the garbage collector could then free the object before the call is actually made.

3.3 Exporting Closures

Life would be rather boring if we could only import foreign functions. Things start to become interesting when we can masquerade Haskell functions as plain function pointers and pass those the outside world.

A static foreign export declaration tells the Haskell compiler to export a declared Haskell function behind a C-callable function interface. For example
(assuming we are on the Win32 platform) we can create a DLL with an entry
**Reverse** that reverses a C string by unmarshalling it, reversing the resulting
Haskell string and then marshalling the reversed string back to a C string as
follows:

```haskell
type CString = Addr

foreign export stdcall "Reverse"
primReverse :: CString -> IO CString

primReverse =
    ( unmarshallString ## return.reverse ## marshallString )
```

The double hash operator `##` is just reverse monadic composition:
\( f ## g = \lambda a \rightarrow \text{do}\{ b <- f a; g b\} \). Later, we will also use the single
hash operator `#`, which is just reverse function application \( a # f = f a \).

Haskell *interpreters*, such as for example Hugs, do not naturally support stati-
cally exported functions. Fortunately, we can also *dynamically* export Haskell
functions as if they are C function pointers\(^3\). Dynamically exported Haskell
functions are extremely powerful and make programming with callbacks very
easy, as the following example demonstrates.

The **JNIEnv** method **RegisterNatives** dynamically registers new native meth-
ods with a running JVM. The **RegisterNatives** entry takes a reference to a
class object in which the native methods will be registered, and an array of
**JNINativeMethod** structures that contain the names, types and implementa-
tions of the native methods:

```c
typedef struct {
    [string] char* name;
    [string] char* signature;
    [ptr] FARPROC fnPtr;
} JNINativeMethod;

jint RegisterNatives
    ( [in] JNIEnv* env
      , [in] jclass clazz
      , [in, size_is(nMethods)] JNINativeMethod* meths
      , [in] jint nMethods
    );
```

To be able to call **RegisterNatives** in Haskell through a **JNIEnv** pointer,
we must first dereference the **JNIEnv** pointer, fetch the 215-th entry in the
function table and coerce that into a Haskell callable function. Then we pass

---

3 In section [4] we will show how static exports can be simulated using dynamic exports
and a statically exported interpreter component.
the resulting function the required arguments to compute the result:

```
getProcAt :: Int -> Addr -> IO FARPROC

foreign import dynamic mkRegisterNatives
  :: FARPROC -> IO (JNIEnv -> Jclass -> Addr -> Int -> IO Int)

registerNatives :: JNIEnv -> Jclass ->
  Addr -> Int -> IO ()
registerNatives = \env -> \clazz -> \meths -> \n -> do{
  f <- (deref
## getProcAt 215
## mkRegisterNatives
) env;
  f env clazz meths n;
}
```

Suppose we want to register a native function with signature
void SayHello (JNIEnv*, jobject) via the `registerNatives` function.
We can easily construct the required function pointer from a Haskell function
with type `JNIEnv -> Jobject -> IO ()` by using function `mkSayHello`.
The latter function is defined using a `foreign export dynamic` declaration:

```
foreign export dynamic
  mkSayHello :: (JNIEnv -> Jobject -> IO ()) -> FARPROC
```

Thus `mkSayHello` is a Haskell function that takes any Haskell function of type
`JNIEnv -> Jobject -> IO ()` and at run-time creates a C function pointer,
that when called will call the exported Haskell function:

```
sayHello_ :: FARPROC
sayHello_ = mkSayHello sayHello
```

Assuming a function `marshallJNINativeMethods` that builds an array of
`JNINativeMethod` structures from a Haskell list of name/signature/function
pointer-triples

```
marshallJNINativeMethods :: [(String,String,FARPROC)] -> IO Addr
```

it is not hard to define a function `registerHello` that registers the
Haskell function `sayHello` as the native implementation of a Java method
`void sayHello()`:

```
registerHello = \env -> \clazz -> do{
  p <- marshallJNINativeMethods [("SayHello","()V", sayHello_)];
  env # registerNatives p 1 clazz;
```
3.4 Interfacing With the GC Using Stable Pointers

When we pass Haskell values (such as dynamically exported functions), to the outside world, we must prevent the Haskell garbage collector from moving them around inside the Haskell heap during a collection. Otherwise, an external function that holds on to the value will suddenly point to a completely different value. In addition, since we have no control over how many copies are made of exported objects, the garbage collector cannot automatically free them anymore. Haskell values are therefore hidden behind an extra level of indirection called stable pointers [5]. For the remainder of this paper we don’t need to understand stable pointers in detail except that we know that they exist behind the scenes.

Now that we have covered the basics of interacting with the outside world and of JNI, we can start looking at the main subject of this paper, the bridge between Java and Haskell via the Java Native Interface (JNI).

4 Calling Java From Haskell and Haskell from Java

Calling Java from Haskell and Haskell from Java is in principle no different from calling Java from C and calling C from Java.

Calling Java from Haskell poses no problems as long as we have static and dynamic foreign import. In order to let Java call Haskell, the only provision is that Java can either load a DLL that contains Haskell implementations of the required native methods (to implement native methods directly), or some other DLL that can dynamically register the native methods using the JRegisterNatives JNI entry we have seen in section 3.3.

The DietHEP component [11] leverages on the foreign function interface we defined for Haskell in previous papers [3] and allows us to view any Haskell module as an ordinary DLL. Clients of DietHEP see no difference (except for the extra flexibility of specifying the calling convention at runtime when using GetProcAddressEx) between using an ordinary DLL via the kernel32.dll or using the DietHEP primitives. Going through the extra level of indirection of DietHEP allows us to abstract from the underlying Haskell implementation (provided of course that it supports the DietHEP interface).

The LoadLibrary function takes the name of a Haskell module (or GHC compiled binary), loads it and returns a handle to the module. Function GetProcAddress takes that handle and a function name and returns a ‘foreign export’-ed version of the requested Haskell function.

[dllname("DietHEP.dll")]}
To register native Haskell methods via DietHEP, we wrap the DietHEP component into a Java class `DietHEP` in the expected way:

```java
class DietHEP {
    static native int LoadLibrary (String n);
    static native int GetProcAddress (int m, String n);
    ...
}
```

We also assume that we have a class `JNI` that reflects (amongst others) the JNI entry `registerNatives` back into Java as the native method `RegisterNatives`:

```java
class NativeMethod {
    String name; String signature; int fnPtr;
    ...
}

class JNI {
    static native void RegisterNatives (Class c, NativeMethod[] ms);
    ...
}
```

Using these two classes, the `Hello` example of section 2.1 can be written to use the Haskell function `Hello.sayHello` by first loading the Haskell module `Hello` that contains the definition of the `sayHello` function, and then registering the dynamically exported function pointer for `sayHello` with the Java `Hello` class:

```java
class Hello {
    public static native void SayHello ();
    static {
        int h = DietHEP.LoadLibrary ("Hello");
        int sayHello = DietHEP.GetProcAddress (h,"sayHello");
        JNI.RegisterNatives ( Hello.class
            , new NativeMethod[]{
            new NativeMethod ("SayHello","()V", sayHello)
```
5 A User Friendly Library

At this stage, we have shown the bare bones of the integration between Haskell and Java. We will make the binding more user friendly in a number of steps. First we will abstract away commonly occurring patterns. Second, we hide the explicit threading of the JNIEnv argument through all JNI related code. Next, we will use overloading to hide the encoding of result-types in JNI methods.

5.1 Combining JNI calls

Calling an instance method void Foo() on a Java object obj and a JNIEnv pointer env is a three step process:

(i) Get the class object for obj using the JNI method GetObjectClass:
   
   cls <- env # getObjectClass obj

(ii) Look up the methodID for the method via its name and type description using the JNI method GetMethodID:

   mid <- env # getMethodID cls "Foo" "()V"

(iii) Call the method using the JNI method CallMethod where <t> is the result type of the method, passing the object, the methodID and the actual arguments:

   env # callVoidMethod obj mid nullAddr

To call a static or instance method, or to get or set a static or instance field, we have to go through the same steps; fetch a class object, look up a method or field ID based on the name and signature, and eventually perform the desired operation. For the sake of presentation however, we will ignore invoking static methods until section 5.8. Space limitations preclude us from explaining (static) field access, but they don’t require new concepts to be dealt with.

Our first step in simplifying interaction with JNI is to combine the basic 3-step JNI call sequence into a family of functions callMethod. These functions take a JObject pointer, the name and signature of the method or field, a ForeignObject pointer to an array of arguments, and return a value of <t> (where as usual, we qualify the primitive callMethod with Prim):

   callMethod :: JObject -> String -> String -> ForeignObject
5.2 Hiding JNIEnv

In previous examples we have seen that most interaction with Java goes via the JNIEnv interface pointer, and that all methods of the JNIEnv interface also take a ‘this’ pointer (or a class object pointer) as an additional argument. Having to thread `env` around is rather tedious, so we are willing to move a few mountains at this moment to save a lot of work later.

The JNIEnv pointer is only valid in its associated thread, so we cannot put it in a global variable just like that. The JavaVM pointer however does remain valid across different threads and we can safely put that in a global variable. To do that, we need to get hold of the JavaVM pointer in which the current thread is running, and we postulate that (current and future) JNI implementations provide some way of doing this. We will follow one of the suggestions of Liang ([8], section 8.1.4) that relies on the fact that the current JVM releases do not support the creation of more than one JVM instance inside a single process ([8], page 254).

The JNI library function `getCreatedJavaVMs` returns a list of all currently running JVMs and as we argued above, we may safely assume that `getCreatedJavaVMs` always returns a singleton list containing the single running JavaVM. Hence `getCreatedJavaVMs` is a pure function and thus we can use `unsafePerformIO` to create a global ‘variable’ `javaVM` that contains the pointer to the running JVM:

```hs
javaVM :: JavaVM
javaVM = unsafePerformIO $ do{
    [javaVM] <- getCreatedJavaVMs;
    return javaVM;
}
```

The JavaVM interface entry `attachCurrentThread` returns the JNIEnv pointer for the current thread. Using global variable `javaVM` we can now easily and safely fetch the currently valid JNIEnv pointer in any context:

```hs
-> JNIEnv -> IO <t>
call<t>Method =
\this -> \name -> \sig -> \a -> \env -> do{
    cls <- env # getObjectClass this;
    mid <- env # getMethodID cls name mid;
    env # Prim.call<t>Method this mid a;
}
```

Our next step will be to eliminate the JNIEnv parameter.

---

The book gives several other suggestions, but this one is the simplest to implement. Another possibility would be to capture the current JVM in the JNI_OnLoad event handler.
attachCurrentThread :: JavaVM -> IO JNIEnv

getJNIEnv :: IO JNIEnv
getJNIEnv = javaVM # attachCurrentThread

Now that we can obtain the right JNIEnv in any context, we don’t have to supply the JNIEnv argument to call<t>Method we defined previously any more, since we can fetch it when we need it (again we qualify the more primitive version of a function with Prim):

\[
\text{call}<t>\text{Method} :: \text{String} \to \text{String} \to \text{ForeignObject} \to \text{Object} \to \text{IO } <t>
\]

\[
\text{call}<t>\text{Method} = \text{\lambda name} \to \text{\lambda sig} \to \text{\lambda args} \to \text{\lambda this} \to \text{do}
\]

\[
\begin{aligned}
\text{env} &\leftarrow \text{getJNIEnv;}
\text{env } &\# \text{Prim.call}<t>\text{Method this name sig args;}
\end{aligned}
\]

5.3 Overloading argument types

Before we start using overloading to simplify JNI, we have to make a short digression on marshalling. In section 3.1, we assumed that we had a function marshallString :: String \to IO CString that takes a Haskell string and returns a pointer to a null terminated array of characters.

The primitive string marshall function marshallString itself is defined in terms two more primitive functions. Function marshallByRefChar :: Char \to [Addr \to IO ()] returns (in this case a singleton) list of functions that writes a character at the indicated addresses. Function writeAt takes a list of Addr \to IO () functions and a start address and calls each function in the list to write the value at subsequent addresses:

\[
\text{writeAt} :: [\text{Addr} \to \text{IO ()}] \to \text{Addr} \to \text{IO ()}
\]

To marshall a string, we allocate enough memory to hold the string, generate a list of marshall functions for each character in the string and then write each one into the allocated memory:

\[
\text{marshallString} :: \text{String} \to \text{IO Addr}
\text{marshallString} = \text{\lambda s} \to \text{do}
\]

\[
\begin{aligned}
\text{let cs} &\leftarrow s++[\text{chr 0}];
\text{p} &\leftarrow \text{malloc } $ \text{length cs};
\text{writeAt} \text{ (concat (map marshallByRefChar cs)} \text{ p;}
\text{return p;}
\end{aligned}
\]

The free function for strings simply calls the system function free :: Addr \to IO (), which deallocates a memory block that was
previously allocated by a call to malloc:

```haskell
freeString :: Addr -> IO ()
freeString = \s -> do{ free s; }
```

In general, we have a family of functions `marshall\(<t>\)`, `marshallByRef\(<t>\)`, `free\(<t>\)` for each \(<t>\) that we can marshall from Haskell to Java:

```haskell
marshall\(<t>\) :: \(<t>\) -> IO Addr
marshallByRef\(<t>\) :: \(<t>\) -> [Addr -> IO ()]
free\(<t>\) :: Addr -> IO ()
```

At this moment bells should start to ring and alarms should go off: “ad-hoc overloading”, this is *exactly* what Haskell type classes were invented for!

So we define a new type class `GoesToJava` with three methods `marshall`, `marshallByRef` and `free`. To be able to overload `free` as well, we add an extra phantom argument to the `free` function in the class:

```haskell
class GoesToJava a where {
    marshall :: a -> IO Addr;
    marshallByRef :: a -> [Addr -> IO ()];
    free :: a -> Addr -> IO ();
}
```

and make all types that can be marshalled from Haskell to Java (Char, Int, Integer, Double, Float, JObject, String, ...) an instance of the `GoesToJava` class:

```haskell
instance GoesToJava \(t\) where {
    marshall = marshall\(<t>\);
    marshallByRef = marshallByRef\(<t>\);
    free = \a -> free\(<t>\);
}
```

The `marshallByRef` instances for Double and Integer return a two element list with the second entry being `\a -> do{ return () }` to ensure that the `marshall` function will allocate enough memory to hold a double or a long.

We continue to overload `GoesToJava` for tuples to encode Java methods with zero, or more than one argument. This is impossible if we use curried functions to represent multiple argument Java methods in Haskell.

```haskell
instance GoesToJava () where {
    marshall = \() -> do{ malloc 0; };
    marshallByRef = \() -> [];
    free = \() -> \p -> do{ free p; };
}
```
instance
  (GoesToJava a, GoesToJava b) =>
  GoesToJava (a,b) where {
    marshall = \(a,b) -> do{
      let ab = marshallByRef a ++ marshallByRef b;
      p <- malloc $ length ab;
      writeAt ab p;
      return p;
    };
    marshallByRef = \(a,b) ->
      marshallByRef a ++ marshallByRef b;
    free = \(a,b) -> \p -> do{
      (deref ## free) p;
      (deref ## free) (incrAddr p);
      free p;
    }
  }

By using the GoesToJava class we don’t have to marshall arguments to a JNI call into an array anymore, but we can simply write (again qualifying the previous version of call<\t>Method with Prim):

call<\t>Method :: GoesToJava a =>
  String -> String -> a -> Jobject -> IO \<\t>
call<\t>Method =
  \name -> \sig -> \a -> \this -> do{
    p <- marshall a;
    r <- newForeignObject p (free a);
    env <- getJNIEnv;
    env # Prim.call<\t>Method this name sig r;
  }

5.4 Overloading result types

The signature of JNI calls above obviously cries out to be overloaded as well by means of a type class ComesFromJava with an instance for each \<\t> ((), JObject, Bool, Byte, Char, Short, Int, Integer, Float, Double) that can be returned by a JNI method invocation:

class ComesFromJava b where {
  callMethod :: GoesToJava a =>
    String -> String -> a -> JObject -> IO b;
}

instance ComesFromJava \<\t> where {
  callMethod = call<\t>Method;
At this point we have reduced the complexity for JNI calls quite significantly when compared with the most primitive form, but there are still ample possibilities to improve. The next thing that we are going to attack is the explicit passing of type descriptor strings, which at the moment is rather unsafe.

5.5 Generating type descriptors

The problem with the current version of function `callMethod` is that there is no connection between the value of the type descriptor string, which denotes the Java type of a field or method, and the corresponding Haskell types of the argument and result of the `callMethod` function. For example, the following definition is accepted by the Haskell type-checker:

```haskell
foo :: () -> JObject -> IO Int
foo = \() -> \obj -> do{ obj # callMethod "foo" "(I)V" (); }
```

Unfortunately, it leads to a runtime error because the `foo` method is called as if it has Java type `void foo(int)` instead of at the correct type `int foo()` (which corresponds to the type descriptor string should be "()I").

We will generate type descriptors by a family of functions `descriptor<\text{t}> :: \text{<t>} \rightarrow \text{String}^5`. As we did in the previous sections, we define a Haskell type class, in this case called `Descriptor`, that witnesses for which \text{<t>} we can obtain a descriptor:

```haskell
class Descriptor a where { descriptor :: a -> String; }

instance Descriptor () where { descriptor = \"\"; }
instance (Descriptor a, Descriptor b) where { descriptor = descriptor a; }
```

The type descriptor for the primitive JNI type `Bool` is "Z":

```haskell
descriptorBool :: Bool -> String
descriptorBool = \z -> "Z"
```

The other basic types are mapped as follows: `Byte` ↦ "B", `Char` ↦ "C", `Short` ↦ "S", `Int` ↦ "I", `Integer` ↦ "L", `Float` ↦ "F", `Double` ↦ "D".

Type descriptors for methods arguments are formed by concatenating the type descriptors of the individual arguments:

```haskell
instance Descriptor () where { descriptor = \"\"; }
instance (Descriptor a, Descriptor b)
```

---

5 This trick is used in the Haskell module `Dynamic` as well, and is attributed to the legendary hacker Lennart Augustsson.
Type descriptors for methods are formed by enclosing the type descriptor for the arguments inside parenthesis followed by the type descriptor of the result type, with the exception that \( V \) is used for the void return type:

```haskell
methodDescriptor :: (Descriptor a, Descriptor b) => a -> b -> String
methodDescriptor = \a -> \b ->
concat
[ "(" , descriptor a ,")"
, if null (descriptor b)
   then "V"
   else descriptor b
]
```

To deploy function `methodDescriptor` we need to jump one more hurdle. In order to build the method descriptor string, we need to know the type of the result of a call. Since the `methodDescriptor` function does not need the value of its arguments to compute a type descriptor, we can recursively build the right method descriptor from a call which will take that descriptor as an argument:

```haskell
callMethod :: ( Descriptor a, GoesToJava a , Descriptor b, ComesFromJava b ) => String -> -> a -> Jobject -> IO b
callMethod = \name -> \a -> \this
-> let{
   sig = methodDescriptor a (unsafePerformIO mb);
   mb = this # Prim.callMethod name sig a
} in mb
```

There are several things to note about the above code. First of all, it intimately relies on lazy evaluation, function `methodDescriptor` only depends on the types of its arguments and not on their values. Secondly, the call to `unsafePerformIO` is there solely to strip the `IO` from the result type of the `callMethod` function, which is completely safe.

### 5.6 Encoding Inheritance using phantom types

The `callMethod` function as we have reached by now is not yet as safe and flexible as it could be. In particular, we do not reflect the fact that in Java classes satisfy an inheritance relation. We can however encode inheritance
using the same trick we used for COM components [6].

For the Java root class `Object` we define a phantom type `Object_` that contains the low level untyped `Jobject` pointer, and immediately hide `Object_` using type synonym `Object` (we will see in just a minute why):

```haskell
data Object_ a = Object Jobject
type Object a = Object_ a
```

Type `Object_` is called a *phantom* type because its type parameter `a` is not used in its right hand side.

For every class `<class>` that extends its superclass `<super>`, we define a new (empty) phantom type `<class>_` and a type synonym `<class>` that expresses that `<class>` extends `<super>`:

```haskell
data <class>_ a
type <class> a = <super> ( <class>_ a )
```

The definition for `<class>` expands to `Object_ (...) ( <super>_ ((<class>_ a)) ...)`, which precisely expresses the inheritance relationship between `Object` and `<super>` and `<class>`.

An instance of the Java class `<class>` is represented by a value of type `<class> ()`, which is still the same old `Jobject` pointer, but now guarded by a little type harness. The nice thing about this is that the Haskell type checker prevents us from using a method of the `<class>` class at an instance for which `<class>` is not a superclass.

Of course, when actually calling a method we use the `Jobject` pointer that is hidden inside the `Object` wrapper:

```haskell
callMethod ::
  (Descriptor a, GoesToJava a , Descriptor b, ComesFromJava b)
  => String -> a -> Object o -> IO b
callMethod = \name -> \a -> \(Object this) ->
  Prim.callMethod name a this
```

5.6.1 Class descriptors

Using the above encoding of Java classes into Haskell types, we can generate type descriptor strings for Java classes by defining a function `descriptor<class>` for each Java class `<class>`:

```haskell
descriptor<class> :: <class> () -> String
descriptor<class> = \_ -> "L<classdescriptor>;;"
```

```haskell
instance Descriptor (<class> ()) where {
  descriptor = descriptor<class>;
}```
A *class descriptor* is the fully qualified Java class with all .-s replaced by /-s; for instance the class descriptor for the Java class `Object` is `java/lang/Object`.

For methods that take objects as arguments, JNI expects the class descriptor for the formal parameter and not of the actual argument. This means for example that if we call the method `Component add(Component comp)` of the Java class `Container` with an argument whose run-time class is `Frame` (a subclass of `Container`), we still must pass the class descriptor `Ljava/awt/Component;` instead of `Ljava/awt/Frame;`. The correct version of `add` silently casts its argument to a `Component ()` just before making the actual call to function `add`:

```haskell
add :: Component a -> Container b -> IO (Component a)
add = \comp -> \this -> do{
  this # callMethod "add" (cast comp :: Component ());
}
```

Inside function `add` we use the function `cast :: a -> b` to change the first argument of type `Component a` (which could be any subclass of `Component`) to exactly `Component ()` so that we get the correct method descriptor. Function `add` itself remains completely type safe.

The above situation is the only case where we have to explicitly up an upcast in Haskell; the question remains how to deal with the occasional need for downcasting. In Java, it is possible to unsafely cast between arbitrary classes, which might raise a dynamic `ClassCastException`. Using the Haskell `cast` to do an illegal arbitrary cast will either result in a `ClassCastException` or other exception being thrown in Java (which is checked using the JNI entry `ExceptionOccured`), or a call to one of underlying the JNI entry calls will signal an error (usually by returning `NULL`). The Haskell wrapper will catch these and propagate the errors to the Haskell level by raising a `userError` when returning. In other words, using `cast` to simulate arbitrary Java casting in Haskell is no more safe or unsafe than arbitrary casting in Java.

### 5.7 Dealing with interfaces

Besides classes, Java also supports *interfaces*. The role of interfaces in Java is quite similar to that of type classes in Haskell. Saying that a Java class implements an interface is like making an instance declaration in Haskell.

To deal with interfaces in Lambada we introduce an empty Haskell type class for each Java interface, possibly by extending some one or more base interfaces, for example:
class ImageObserver a
class MenuContainer a
class Serializable a
class EventListener a
class EventListener a => WindowListener a

For each method in an interface, we generate a function that requires its argument to satisfy the class constraint, eg for the MenuContainer interface we define:

```haskell
getFont :: MenuContainer a => () -> a -> IO (Font ())
getFont = \\() -> \\this -> do{
  this # callMethod "getFont" ()
}
```

For each Java class `<class>` that implements an interface `<interface>`, we make the corresponding Haskell type `<class>` an instance of the Haskell type class `<interface>` that corresponds to the Java interface `<interface>`. For instance, the Java class Component implements the interfaces ImageObserver, MenuContainer, and Serializable, so in Haskell we write:

```haskell
instance ImageObserver (Component a)
instance MenuContainer (Component a)
instance Serializable (Component a)
```

Note that we have to be careful not to write Component (), since every subclass of Component implements the interfaces that Component implements.

Since Component is an instance of MenuContainer, a call `c # getFont ()` is well-typed whenever `c` has type Component (), Container (), Window (), Frame (), or any subclass of these.

Some methods expect interface ‘instances’ as arguments:

```haskell
void addWindowListener(WindowListener)
```

To deal with this situation, we have to be able to generate descriptor strings for interfaces. For each interface `<interface>` we define a new type `<interface>` that uniquely identifies that interface and make `<interface>` an instance of Descriptor. Of course an interface satisfies its own interface constraint (in reality, we have to do name mangling because in Haskell data types and type classes are in the same name spaces but here no confusion can arise):

```haskell
data <interface>

descriptor<interface> :: <interface> -> String
descriptor<interface> = \\_ -> "L<interface>";
```
instance Descriptor <interface> where {
    descriptor = descriptor<interface>;
}

instance interface <interface>

Now we can (safely) coerce any object reference that implements a certain interface to that interface and obtain the correct descriptor.

5.8 Static methods and class objects

Up to now we have ignored static methods. A static or class method is a method that does not act on a class instance, but on the class object instance.

Calling a static method, say void Foo(), on a class object cls and given a JNIEnv pointer env is quite similar to calling an instance method, except that the method is not invoked on the object instance but on its class object:

(i) Look up the methodID for the static method via its name and type description using the JNI method GetStaticMethodID:

   mid <- env # getStaticMethodID cls "Foo" "()V"

(ii) Call the static method using the JNI method CallStaticMethod where <t> is the result type of the method, passing the class object, the methodID and the actual arguments:

   env # callStaticVoidMethod cls mid nullAddr

If we continue and mirror the overloading have done for instance methods, we will end up with a function callStaticMethod for invoking static methods on a class object:

   callStaticMethod ::
   (Descriptor a, GoesToJava a , Descriptor b, ComesFromJava b) => String -> a -> jclass -> IO b

Although in Java we can call a static method on an object (which the compiler translates to the right class), the recommended style is to call it on the class directly. By requiring the latter style, we can distinguish the types of callStaticMethod and callMethod. This enables us to introduce overloading to hide the distinction between calling static and instance methods:

   class JNI j where {
   call
     :: ( Descriptor a, GoesToJava a
      , Descriptor b, ComesFromJava b
     ) => String -> a -> j -> IO b;
   }
For instance method calls we can immediately call `callMethod`:

```haskell
instance JNI (Object a) where { call = callMethod; }
```

For static methods we define a new phantom type `ClassObject` such that `ClassObject <class>` uniquely encodes the class object of a Java class `<class>`. Type `ClassObject` just wraps a type harness around an untyped `Jclass` pointer:

```haskell
data ClassObject a = ClassObject Jclass
```

The JNI instance declaration for static methods peels off the type harness and calls `callStaticMethod`:

```haskell
instance JNI (ClassObject a) where {
  call = 
    \n -> \a -> (ClassObject clazz) ->
      clazz # callStaticMethod n a;
}
```

So now we can call static methods on class objects, but how do we obtain class objects in the first place? The `JNIEnv` entry `findClass` takes a descriptor for the Java class `<class>` and returns the class object for `<class>`:

```haskell
getClassObject :: Descriptor (Object o) => (Object o) -> ClassObject (Object o)
getClassObject = \this ->
  unsafePerformIO $ do{
    env <- getEnv;
    clazz <- env # findClass (descriptor this);
    return (ClassObject clazz);
  }
```

Since function `getClassObject` does not use its argument to compute the class object, we can define an overload class object ‘constant’ with the help of function `stripObject :: Object o -> o`:

```haskell
classObject :: Descriptor (Object o) => ClassObject (Object o)
classObject = let{ c = getClassObject (stripObject c) } in c
```

### 5.9 Constructing object instances

The distinction between instances (subclasses of `java.lang.Object`) and class objects (subclasses of `java.lang.Class`) can be confusing at first, especially because we can also call non-static methods on class references. One such example are constructor functions to create class instances.

Without using overloading to generate method signatures, we can combine the steps to create a new object into a single function as follows:
new :: (GoesToJava a) => String -> Jclass -> a -> IO (Object o)
new = \sig -> \clazz -> \a -> do{
cid <- env # getMethodId clazz "<init>" sig;
p <- marshall a;
r <- newForeignObject (free a) p;
o <- env # newObject clazz cid r;
return (Object o);
}

As we did before, we want to use function methodDescriptor to generate the
method descriptor, and furthermore we want to avoid passing the class object
as an argument to new since we can get that using function getClassObject
(again we qualify the previous version of new with Prim):

new :: (GoesToJava a , Descriptor (Object o) ) => a -> IO (Object o)
new = \a -> let{
o = unsafePerformIO mo;
clazz = getClassObject o;
sig = methodDescriptor a o;
mo = Prim.new sig clazz a;
} in mo

In practice overloading new on its return type often requires manual resolution
using explicit type signatures. For convenience we will therefor define an
additional un-overloaded constructor function new_<class> for each <class>
that is an instance of Descriptor by restricting the type of new:

new_<class> :: (GoesToJava a) => a -> IO <class>
new_<class> = new

6 Simulating anonymous inner classes

Haskell functions/actions are represented in Java as objects of class Function,
which store the actual exported Haskell function pointer as an int in a private
field f and implement a native method call that takes an Object array as
argument, and an Object representing the this pointer (or class object) and
returns an Object as result:

class Function {
    private int f;
    Function(int f){ this.f = f; }
    public native Object call
        (Object[] args, Object this);
    ...
}

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A typical situation is that we have to create an event-handler that implements the WindowListener interface. To facilitate the construction of such a handler in Haskell, we define a Java class WindowHandler that has a private Function object field for each of its methods which are initialized by the constructor function. Each method in the WindowHandler class simply delegates its work to the appropriate Function object:

```java
class WindowHandler
    implements WindowListener {
    private Function activated;
    ...
    WindowHandler
        (Function activated, ...) {
            this.activated = activated;
            ...
        }
        void windowActivated(WindowEvent e){
            this.activated.call (new Object [] {e}, this);
        }
        ...
    }
```

The JNI framework dictates that the native implementation of the call method is a function that takes the JNIEnv and the JObject this-pointer as all native methods do, a JObject argument that represents the argument array, the explicit this argument of the instance in which the call was made, and that returns a JObject:

```haskell```
call :: JNIEnv -> JObject -> (JObject -> JObject -> IO JObject)
```

Function call uses the this pointer of the Function instance to retrieve the function pointer from the private field f and then calls that as a function of Haskell signature (JObject -> JObject -> IO JObject):

```haskell```
foreign import dynamic importFunction
    :: FARPROC
    -> IO (JObject -> JObject -> IO JObject)
```

```haskell```
call = \env -> \this -> \args -> \t -> do{
    f <- env # getField this "f";
    importFunction f args t;
}
```

We create Function instances in Haskell by creating a C function pointer for the Haskell function we want to export using exportFunction and then use the Haskell JNI library to call the new_Function constructor:
foreign export dynamic exportFunction
    :: (Jobject -> Jobject -> IO Jobject) -> FARPROC
new_Function :: Int -> IO (Function ()

All this quite low-level and error prone, so let’s use a type class to tidy up and
do the grungy work for us.

First of all, we extend the ComesFromJava class with a new member func-
tion unmarshallArray that unmarshalls zero or more values from a Jobject
pointer representing an Object array to a tuple of corresponding Haskell val-
ues, and a new member function unmarshall that marshalls a single Jobject
pointer into its corresponding Haskell value:

class ComesFromJava b where {
    callMethod :: GoesToJava a =>
        String -> String -> a -> Object o -> IO b;
    unmarshall :: Jobject -> IO b;
    unmarshallArray :: Jarray -> IO b
}

For basic types <t> we simply use the unmarshaller for that type, or first get
the single element from the array and unmarshall that:

instance ComesFromJava <t> where {
    unmarshall = \o -> do{
        unmarshall<t> o;
        unmarshallArray = \args
            args # (getArrayElement 0 ## unmarshall);
    }
}

For pairs, triples, etc, we recursively unmarshall the elements from the array
one at a time.

instance
    (ComesFromJava a, ComesFromJava b)
=> ComesFromJava (a,b) where {
    unmarshall = unmarshallArray;
    unmarshallArray = \args -> do{
        a <- args # (getArrayElement 0 ## unmarshall);
        b <- args # (getArrayElement 1 ## unmarshall);
        return (a,b);
    }
}

Function marshalFun takes a high-level Haskell function of type
(ComesFromJava a, GoesToJava b) => (a -> Object o -> IO b), turns
it into a low-level function of type JObject -> JObject -> IO JObject and
then dynamically exports it using exportFunction:
marshallFun
  :: (ComesFromJava a, GoesToJava b)
  => (a -> Object o -> IO b) -> IO FARPROC
marshallFun = \f -> exportFunction $ \args -> \this -> do{
  a <- unmarshall args;
  t <- unmarshallObject this
  b <- f a t; marshall b;
}

With the help of marshallFun we can finally make functions an instance
of GoesToJava. To marshall a function from Haskell to Java, we first use
marshallFun to get its function pointer representation, then we construct a
new Java Function object for the pointer and marshall the resulting Java
object to obtain a pointer we can throw over to Java:

instance
  (ComesFromJava a, GoesToJava b)
  => GoesToJava (a -> Object o -> IO b)
  where {
    marshall = \f -> do{
      a <- marshallFun f;
      func <- new_Function (toInt a);
      marshall func
    }
  }

Since functions are an instance of GoesToJava, writing event handlers in
Java has now become completely painless. We can call the constructor
for our WindowHandler example, passing it seven Haskell functions of type
Event e -> WindowHandler w -> IO () that handle each of the seven pos-
sible Window events:

new_WindowHandler
  ( \event -> \this -> do{
      this # activated event;
    }, ... )

This nearly exactly matches the corresponding code we would write in Java

new WindowHandler () { void activated (Event e) { ...; }... }

7 Tool support

To reiterate the introduction of this paper, Java is an interesting partner for
a foreign function interface because there are a lot of Java classes out there.
However, even with the layers of abstractions we have introduced on top of
JNI, manually coding the Haskell callable stubs for the Java APIs we are interested in is not always practical.

Lambda provides tool support in concert with HDirect. HDirect is IDL based, so provided we can derive an IDL specification from a Java class or interface, HDirect will take care of generating valid Haskell stubs. We do this by providing a tool that uses the Java Reflection API to generate IDL interfaces (annotated with a couple of custom attributes) corresponding to Java classes or interfaces. Provided with that information, HDirect is then able to automatically generate the Haskell callable stubs we have described in this paper. HDirect is also capable of generating Java callable stubs to Haskell implemented classes/methods.

At the Lambda implementation level, HDirect was also used to generate the low-level JNI stubs outlined in Section 3.

8 Related work

Claus Reinke was a remarkably early adaptor of JNI, as early as 1997 he proposed to use JNI to interoperate between Haskell and Java. Reinke’s system provides only one-way interop between Haskell and Java using the basic low-level JNI methods, i.e., he does not support calling back from Java to Haskell. The most important contribution of our work over his is our sophisticated use of overloading and other advanced language features to hide the idiosyncrasies of JNI. The integration that Lambda offers between Haskell and Java is very close to the tight integration offered by MLj, with the important exception that Lambda is implemented completely within Haskell, whereas MLj needs many language changes.

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