The Capacity and Optimal Signaling for Gaussian MIMO Channels Under Interference Constraints

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Abstract—Gaussian MIMO channel under the joint transmit (Tx) and interference power constraints (TPC and IPC) is studied. A closed-form solution for the optimal Tx covariance matrix in the general case is obtained using the KKT-based approach, up to dual variables. A number of more explicit closed-form solutions are given with optimal dual variables, including full-rank and rank-1 (beamforming) cases as well as the case of identical eigenvectors (typical for massive MIMO settings), which differ significantly from the standard water-filling solution. A “whitening” filter is shown to be an important part of optimal precoding under interference constraints. Sufficient and necessary conditions for each constraint to be redundant are given. Capacity scaling with the SNR is shown to be determined by a natural linear-algebraic structure of sub-spaces induced by channel matrices of multiple users. A number of unusual properties of optimal Tx covariance matrix under the joint TPC and IPC are pointed out and a bound on its rank is established. An interplay between the TPC and IPC is studied, including the transition from power-limited to interference-limited regimes as the Tx power increases. While closed-from solutions for optimal dual variables are given in some special cases, an iterative bisection algorithm (IBA) is proposed to find optimal dual variables in the general case and its convergence is proved for some special cases. Numerical experiments illustrate its efficient performance. Bounds for the optimal dual variables are given.

Index Terms—MIMO, channel capacity, interference constraint, optimal signaling.

I. INTRODUCTION

EXPONENTIALLY-GROWING volume of high-rate mobile wireless traffic stimulated active development of 5G standards and systems. Due to very high expectations, several new key technologies have been identified to meet those demands, including massive MIMO, millimeter waves (mmWave) and non-orthogonal multiple-access [1]. The ultimate goal is to increase significantly the available bandwidth as well as spectral efficiency to meet the growing traffic demands. However, aggressive frequency re-use and non-orthogonal access schemes can potentially generate significant amount of inter-user interference, which thus has to be carefully managed [2]–[5]. This is somewhat similar to cognitive radio (CR) systems, which also emerged as a powerful approach to exploit underutilized spectrum and hence possibly resolve the spectrum scarcity problem [6]. In both settings, allowing spectrum re-use calls for a careful management of interference. In this respect, multi-antenna (MIMO) systems have high potential due to their significant signal processing capabilities, including interference cancellation and precoding [7], which can also be done in an adaptive and distributed manner [13]. A promising approach is to limit interference to primary receivers (PR) by properly designing secondary transmitters (Tx) while exploiting their multi-antenna capabilities.

The capacity and optimal signalling for the Gaussian MIMO channel under the total Tx power constraint (TPC) is well-known: the optimal (capacity-achieving) signaling is Gaussian and on the eigenvectors of the channel with power allocation to the eigenmodes given by the water-filling (WF) [7]–[9]. Under the more practical per-antenna power constraints (PAC), in addition or instead of the TPC, Gaussian signalling is still optimal but not on the channel eigenvectors anymore so that the standard water-filling solution over the channel eigenmodes does not apply [10], [11]. Much less is known under the added interference power constraint (IPC), which limits the power of interference induced by a secondary transmitter to a primary receiver. A game-theoretic approach to this problem was proposed in [13], where a fixed-point equation was formulated from which the optimal covariance matrix can in principle be determined. Unfortunately, no closed-form solution is known for this equation. In addition, this approach is limited in the following respects: the channel to the primary receiver is required to be full-rank hence excluding the important case of single-antenna devices communicating to a multi-antenna base station or, in general, the cases where the number of Rx antennas is less than the number of Tx antennas - typical for massive MIMO downlink; the TPC is not included explicitly (rather, being “absorbed” into the IPC), hence eliminating the important case of inactive IPC (since this is the only explicit constraint); consequently, no interplay between the TPC and the IPC can be studied.

Earlier studies on cognitive radio MIMO system optimization under interference constraint using game-theoretic approach are extended to the case of channel uncertainty in [14] by developing a number of numerical algorithms for Tx optimization under the global (total) interference power constraint. Due to the non-convex nature of the original problem, a number of approximate and sub-optimal approaches are
adopted, for which provable convergence to a global optimum is out of reach. No closed-form solutions are known for this setting either. Weighted sum-rate maximization in multiuser MISO channel is considered in [15] under interference constraints and numerically-efficient algorithms for Tx optimization based on zero-forcing beamforming are developed, but no closed-form solutions are obtained for this problem. Gaussian MIMO broadcast (BC) and multiple-access channels (MAC) are considered in [16] under general linear Tx covariance constraint, which can also be interpreted as interference constraint, and the earlier BC-MAC duality result is extended to this more general setting. However, no closed-form solutions are obtained for an optimal Tx covariance matrix. In the related context of physical-layer security, the secrecy capacity of the Gaussian MIMO wiretap channel under interference constraints has been characterized in [12] as a non-convex maximization problem (over feasible Tx covariance matrices under the TPC and IPC) or as a convex-concave max-min problem (where “min” is over noise covariance matrices of a genie-aided channel), for which no closed-form solution is known in general.

Unlike most of the studies above, we concentrate here on analysis of the problem and obtain a number of closed-form solutions, which are validated via numerical experiments. This provides deeper understanding of the problem and a number of insights unavailable from numerical algorithms alone. The key difficulty is due to the fact that the feasible set under the IPC is not isotropic anymore (except for some trivial cases) and hence the standard analytical tools (e.g. Hadamard inequality) cannot be used. We overcome this difficulty by using the KKT-based approach and obtain a number of closed-form solutions for an optimal Tx covariance matrix under the TPC and IPC (in Sections III, IV, VI). Both constraints are included explicitly and hence anyone is allowed to be inactive. This allows us to study the interplay between the power and interference constraints and, in particular, the transition from power-limited to interference-limited regimes as the Tx power increases. As an added benefit, no limitations is placed on the rank of the channel to the PR, so that the number of antennas of the PR can be any. Under the added IPC, independent signaling is shown to be sub-optimal for parallel channels to the intended receiver (Rx), unless the PR channels are also parallel or if the IPC is inactive. These results are further extended to multi-user setting with multiple IPCs (in Section IX).

Optimal signaling for the Gaussian MIMO channel under the TPC and the IPC has been also studied in [17], [18] using the dual problem approach, and was later extended to multi-user settings in [19]. However, constraint matrices are required to be full-rank and no closed-form solution was obtained for optimal dual variables. Hence, various numerical algorithms or sub-optimal solutions were proposed (e.g. partial channel projection). Similar problem has been also considered in [20] under multiple linear constraints at the transmitter and an iterative numerical algorithm was developed via a min-max reformulation of the problem. However, no closed-form solution was obtained and the properties of optimal signaling as well as those of the capacity were not studied.

Here, a closed-form solution for an optimal Tx covariance matrix is obtained in the general case (up to dual variables), its properties are studied and a number of more explicit solutions are obtained in some special cases (including explicit solutions for optimal dual variables). Our KKT-based approach does not require full-rank constraint matrices and includes explicit equations for the optimal dual variables, which can be solved efficiently. To this end, we propose an iterative (gradient-free) bisection algorithm (IBA) and prove its convergence (in Sec. VIII). Numerical experiments demonstrate its efficient performance (in Sec. X). Bounds to the optimal dual variables are derived, which facilitate numerical solutions. Properties of the optimal Tx covariance as a function of dual variables are explored: the total Tx power as well as interference power are shown to be decreasing functions of dual variables, which is an important part in the proof of the IBA convergence. In some cases, our KKT-based approach leads to explicit closed-form solutions for the optimal dual variables, including full-rank and rank-1 (beamforming) solutions and the conditions for their optimality. A “whitening” filter is shown to play a prominent role in optimal precoding under interference constraints. Partial null forming, well-known in the antenna array literature [36], is shown to be optimal from the information-theoretic perspective as well in certain cases.

The presented solutions for an optimal Tx covariance matrix under interference constraints and the respective capacity possess a number of unusual properties (not found in the standard WF procedure), namely: an optimal covariance is not necessarily unique; its rank can exceed the main channel rank; the TPC can be inactive; signaling on the main channel eigenmodes is not optimal anymore; the capacity can be zero for non-zero channel and Tx power, and it can be bounded even if the Tx power grows unbounded. Ultimately, these are due to an interplay between the TPC and the IPC.

A simple rank condition is given to characterize the cases where spectrum sharing is possible for any interference power constraint. In general, the primary user has a major impact on the capacity at high SNR while being negligible at low SNR. The capacity scaling with the Tx power under multiple IPCs (primary users) can be understood in terms of a natural linear-algebraic structure of the sub-spaces induced by MIMO channel matrices of multiple users.

The presented closed-form solutions of optimal signaling can be used directly in massive MIMO settings. Since numerical complexity of generic convex solvers can be prohibitively large for massive MIMO (in general, it scales as $m^6$ with the number $m$ of antennas), the above analytical solutions are a valuable low-complexity alternative.

In all considered cases, optimal Tx covariance matrices are significantly different from those of the standard Gaussian MIMO channel under the TPC and/or the PAC [10], [11], and from those of the wiretap channel in [12]. In the latter two cases, optimal Tx covariance matrix remains unknown in the general case while some special cases have been solved.

Finally, it should be pointed out that the channel model we consider here (the standard point-to-point Gaussian MIMO channel under interference constraints at the transmitter) is different from the Gaussian interference channel (G-IC) where multi-user interference is present at each receiver and
there is no interference constraint at the transmitters, as in e.g. [21]–[29]. In the latter case, the capacity and optimal signaling are not known in general, even for the 2-user SISO channel [21] (it is not even known whether Gaussian signaling is optimal in general), except for some special cases, such as strong and weak interference regimes [22], [23], so that various bounds [24], [25] and ad-hoc signaling techniques [26]–[29] are used instead. On the contrary, optimal signaling is known to be Gaussian for the channel model considered here and the capacity can be expressed as an optimization problem over all feasible transmit covariance matrices. Analytical solutions of this problem in the general as well as some special cases and their properties are the main contributions of the present paper.

Notations: bold capitals (\(R\)) denote matrices while bold lower-case letters (\(x\)) denote column vectors; \(R^+\) is the Hermitian conjugation of \(R\); \(R \succeq 0\) means that \(R\) is positive semi-definite; \(|R|\), \(\text{tr}(R)\), \(r(R)\) denote determinant, trace and rank of \(R\), respectively; \(\lambda_i(R)\) is \(i\)-th eigenvalue of \(R\); unless indicated otherwise, eigenvalues are in decreasing order, \(\lambda_1 \geq \lambda_2 \geq \ldots : [\cdot]^{+}\) denotes ceiling, while \((x)^+ = \text{max}[0, x]\) is the positive part of \(x\); \(\mathcal{R}(R)\) and \(\mathcal{N}(R)\) denote the range and null space of \(R\) while \(R^\dagger\) is its Moore-Penrose pseudo-inverse; \(\mathbb{E}\{\cdot\}\) is statistical expectation.

II. CHANNEL MODEL

Let us consider the standard discrete-time model of the Gaussian MIMO channel:

\[ y_1 = H_1 x + \xi_1 \]  

(1)

where \(y_1, x, \xi_1\) and \(H_1\) are the received and transmitted signals, noise and channel matrix, of dimensionality \(n_1 \times 1, m \times 1, n_1 \times 1,\) and \(n_1 \times m,\) respectively, where \(n_1, m\) are the number of Rx and Tx antennas. This is illustrated in Fig. 1. The noise is assumed to be complex Gaussian with zero mean and unit variance, so that the SNR equals to the signal power. A complex-valued channel model is assumed throughout the paper, with full channel state information available both at the transmitter and the receiver. Gaussian signaling is known to be optimal in this setting [7]–[9] so that finding the channel capacity \(C\) amounts to finding an optimal transmit covariance matrix which can be expressed as the following optimization problem (P1):

\[ (P1) : C = \max_{R \in S_0} C(R) \]  

(2)

where \(C(R) = \log |I + W_1 R|\) is an achievable rate with the covariance matrix \(R, W_1 = H_1^H H_1\) is the channel Gramm matrix, \(R\) is the Tx covariance matrix and \(S_0\) is the constraint set. In the case of the total Tx power constraint (TPC) only, it takes the form

\[ S_R = \{ R : R \succeq 0, \text{tr}(R) \leq P_T \}, \]  

(3)

where \(P_T\) is the maximum total Tx power. The solution to this problem is well-known: optimal signaling is on the eigenmodes of \(W_1\), so that they are also the eigenmodes of optimal covariance \(R^\dagger\), and the optimal power allocation is via the water-filling (WF). This solution can be compactly expressed as follows:

\[ R^* = R_{WF} = (\mu^{-1} I - W_1^{-1})_+ \]  

(4)

where \(\mu \geq 0\) is the “water” level found from the total power constraint \(\text{tr}(R^*) = P_T\) and \((R)_+\) denotes positive eigenmodes of Hermitian matrix \(R\):

\[ (R)_+ = \sum_{i, \lambda_i > 0} \lambda_i u_i u_i^+ \]  

(5)

where \(\lambda_i, u_i\) are \(i\)-th eigenvalue and eigenvector of \(R\). Note that this definition allows \(W_1\) in (4) to be singular, since the respective summation (as in (5)) includes only strictly-positive modes: \(R_{WF} = \sum_{i, \lambda_i > 0} (\mu^{-1} - \lambda_i^{-1}) u_i u_i^+\), where \(\lambda_i, u_i\) are \(i\)-th eigenvalue and eigenvector of \(W_1\), so that \(\lambda_i = 0\) are excluded from the summation due to \(\lambda_i \geq \mu \geq 0\).

In a multi-user system, there is a 2nd channel from the Tx to an external (primary) user \(U\), see Fig. 1,

\[ y_2 = H_2 x + \xi_2 \]  

(6)

where \(H_2\) is the matrix of the Tx-U channel (see Fig. 1); all vector/matrix dimensions are equal to the respective number of antennas. There is a limit on how much interference the Tx can induce (via \(x\)) to the user \(U\):

\[ \mathbb{E}\{x^+ H_2^+ H_2 x\} = \text{tr}(H_2^* R H_2^+) \leq P_I \]  

(7)

where \(P_I\) is the maximum acceptable interference power and the left-hand side is the actual interference power at user \(U\). In this setting, the constraint set becomes

\[ S_R = \{ R : R \succeq 0, \text{tr}(R) \leq P_T, \text{tr}(W_2 R) \leq P_I \}, \]  

(8)

where \(W_2 = H_2^+ H_2^+\); this is extended to multiple IPCs in Section IX. The IPC in (8) is consistent with the respective per-user constraints in the CR setting in [13]–[17], with the relay system setting in [18] as well as with the wiretap channel setting in [12]. The Gaussian signalling is still optimal in this setting and the capacity subject to the TPC and IPC can still be expressed as in (2) but the optimal covariance is not \(R_{WF}\) anymore, as discussed in the next section.

III. OPTIMAL SIGNALLING UNDER INTERFERENCE CONSTRAINT

The following Theorem gives a closed-form solution for the optimal Tx covariance matrix under the TPC and the IPC in (8) in the general case.

Theorem 1: Consider the capacity of the Gaussian MIMO channel in (2) under the joint TPC and IPC in (8),

\[ C = \max_{R \in S_0} C(R) \text{ s.t. } R \succeq 0, \text{tr}(R) \leq P_T, \text{tr}(W_2 R) \leq P_I \]  

(9)
The optimal Tx covariance matrix to achieve the capacity can be expressed as follows:

$$R^* = W^\dagger_\mu (I - W^\dagger_\mu W^\dagger_\mu I W^\dagger_\mu) + W^\dagger_\mu$$

(10)

where $W^\dagger_\mu = \{\mu_1 I + \mu_2 W^\dagger_2\}^{-1} W^\dagger_\mu$ is Moore-Penrose pseudo-inverse of $W^\mu$; $\mu_1, \mu_2 \geq 0$ are Lagrange multipliers (dual variables) responsible for the total Tx and interference power constraints found as solutions of the following non-linear equations:

$$\mu_1(tr(R^*) - P_T) = 0, \mu_2(tr(W^2 R^*) - P_T) = 0$$

(11)

subject to $tr(R^*) \leq P_T$, $tr(W^2 R^*) \leq P_I$. The capacity can be expressed as follows:

$$C = \sum_{i: \lambda_{ai} > 1} \log \lambda_{ai}$$

(12)

where $\lambda_{ai} = \lambda_i(W^\dagger_i W^\dagger_1 W^\dagger_\mu)$.

Proof: See Appendix.

This solution can be further extended to multiple IPCs - see Section IX. The known special cases follow from (10): when $W^\mu$ is full-rank, $W^\dagger_1 = W^\dagger_\mu$ [31] and $R^*$ in (10) reduces to the respective solutions in [17], [18]; if the IPC is inactive, then $\mu_2 = 0$, $W^\mu = \sqrt{\mu_2} I$ and $R^* = R_{WF}$, as it should be. Based on (10), one observes that $W^\mu$ plays a role of a preceding “whitening” filter (at the transmitter), which disappears when the IPC is inactive. Note also that, unlike the standard solution under the TPC alone, the optimal signaling in (10) is not on the eigenmodes of $W^1$, unless $W^1$ and $W^2$ have the same eigenvectors or the IPC is inactive ($\mu_2 = 0$). In particular, independent signaling (diagonal $R$) is not optimal even if $W^1$ is diagonal, unless $W^2$ is diagonal as well or $\mu_2 = 0$.

Next, we explore some general properties of the capacity. It is well-known that, without the IPC, $C(P_T)$ grows unbounded as $P_T$ increases, $C(P_T) \rightarrow \infty$ as $P_T \rightarrow \infty$ (assuming $W^1 \neq 0$). This, however, is not necessary under the IPC with fixed $P_I$. The following proposition gives sufficient and necessary conditions when it is indeed the case.

**Proposition 1:** Let $0 \leq P_I < \infty$ be fixed. Then, the capacity grows unbounded as $P_T$ increases, i.e., $C(P_T) \rightarrow \infty$ as $P_T \rightarrow \infty$, if and only if $N(W^2) \notin N(W^1)$.

Proof: See Appendix.

Since the condition of this Proposition is both sufficient and necessary for the unbounded growth of the capacity, it gives the exhaustive characterization of all the cases where such growth is possible. In practical terms, those cases represent the scenarios where any high spectral efficiency is achievable given enough power budget. On the other hand, if $N(W^2) \in N(W^1)$, then the capacity grows unbounded with $P_T$ under the latter condition.

In the standard Gaussian MIMO channel without the IPC, $C = 0$ if either $P_T = 0$ or $W^1 = 0$, i.e., in a trivial way. On the other hand, in the same channel under the TPC and IPC, the capacity can be zero in non-trivial ways, as the following proposition shows. In practical terms, this characterizes the cases where interference constraints of primary users rule out any positive rate of a given user and, hence, spectrum sharing is not possible.

**Proposition 2:** Consider the Gaussian MIMO channel under the TPC and IPC and let $P_T > 0$. Its capacity is zero if and only if $P_I = 0$ and $N(W^2) \in N(W^1)$.

Proof: To prove the “if” part, observe that $tr(W^2 R) = P_I = 0$ implies that $W^2 R = 0$ (since $W^2 R$ has positive eigenvalues, $\lambda_i(W^2 R) \geq 0$) so that $R(R) \notin N(W^2) \in N(W^1)$ and hence $W^2 R = 0$ so that $\log |I + W^1 R| > 0$ for any feasible $R$. Hence, $C = 0$.

To prove the “only if” part, assume first that $P_I > 0$ and set $R = p I$, where $p = \min\{P_T, P_I/(\mu_1(W^1))\}$. Note that $R$ is feasible: $tr(R) \leq P_T$ and $tr(W^2 R) \leq P_I$. Furthermore,

$$C \geq \log |I + W^1 R| > 0$$

(13)

and hence $P_I = 0$ is necessary for $C = 0$. To show that $N(W^2) \notin N(W^1)$ is necessary as well, assume that $N(W^2) \notin N(W^1)$, which implies that $\exists u : W^2 u = 0$, $W^2 u \neq 0$. Now set $R = P_I u u^T$, for which $tr(R) = P_I$, $tr(W^2 R) = 0$, so it is feasible and

$$C \geq \log |I + W^1 R| = \log(1 + P_T u u^T W^1 u) > 0$$

(14)

so that $N(W^2) \notin N(W^1)$ is necessary for $C = 0$.

Note that the condition $P_I = 0$ is equivalent to zero-forcing transmission, i.e., the capacity is zero only if the ZF transmission is required; otherwise, $C > 0$. The condition $N(W^2) \notin N(W^1)$ cannot be satisfied if $r(W^1) > r(W^2)$ and hence $C > 0$ under the latter condition, which is also sufficient for unbounded growth of the capacity with $P_T$. This is summarized below.

**Corollary 1:** If $r(W^1) > r(W^2)$, then
1. $C \neq 0 \forall P_T > 0$.
2. $C(P_T) \rightarrow \infty$ as $P_T \rightarrow \infty \forall P_I > 0$

Thus, the condition $r(W^1) > r(W^2)$ represents favorable scenarios where spectrum sharing is possible for any $P_I$ and arbitrary large capacity can be attained given enough power budget.

It should be pointed out that (11) in Theorem 1 allows anyone of the dual variables to be inactive (i.e. $\mu_1 = 0$ or $\mu_2 = 0$, but not simultaneously), unlike the standard WF solution, where the TPC is always active. While it is not feasible to find dual variables $\mu_1, \mu_2$ in a closed form in general (since (11) is a system of coupled non-linear equations), they can be found in such form in some special cases, as the next sections show. Section VIII will develop an iterative bisection algorithm (IBA) to find the optimal dual variables in the general case with any desired accuracy and prove its convergence.

**IV. FULL-RANK SOLUTIONS**

While Theorem 1 establishes a closed-form solution for optimal covariance $R^*$ in the general case, it is expressed via dual variables $\mu_1, \mu_2$ for which no closed-form solution is known in general so they have to be found numerically.
using (11). This limits insights significantly. In this section, we explore the cases when the optimal covariance $\mathbf{R}^*$ is of full rank and obtain respective closed-form solutions. First, we consider a transmit power-limited regime, where the IPC is redundant and hence the TPC is active.

**Proposition 3:** Let $\mathbf{W}_1 > 0$ and $P_T$ be bounded as follows:

$$m \lambda_1(\mathbf{W}_1^{-1}) - \text{tr}(\mathbf{W}_1^{-1}) < P_T$$

$$\leq \frac{m}{\text{tr}(\mathbf{W}_2)}(P_T + \text{tr}(\mathbf{W}_2 \mathbf{W}_1^{-1})) - \text{tr}(\mathbf{W}_1^{-1})$$

then $\mu_2 = 0$, i.e. the IPC is redundant; $\mathbf{R}^*$ is of full-rank and is given by:

$$\mathbf{R}^* = \mu_1^{-1} \mathbf{I} - \mathbf{W}_1^{-1}$$

(16)

where $\mu_1^{-1} = m^{-1}(P_T + \text{tr}(\mathbf{W}_1^{-1}))$. The capacity can be expressed as

$$C = m \log \frac{P_T + \text{tr}(\mathbf{W}_1^{-1})}{m} + \log |\mathbf{W}_1|$$

(17)

**Proof:** It follows from (10) that, if $\mathbf{R}^*$ is full rank, then so is $\mathbf{W}_\mu$ and

$$(\mathbf{I} - \mathbf{W}_\mu \mathbf{W}_1^{-1} \mathbf{W}_\mu)_{+} = \mathbf{I} - \mathbf{W}_\mu \mathbf{W}_1^{-1} \mathbf{W}_\mu > 0$$

(18)

(since it is of full rank) and hence

$$\mathbf{R}^* = \mathbf{W}_\mu^{-2} - \mathbf{W}_1^{-1}$$

(19)

The full-rank condition $\mathbf{R}^* > 0$ is equivalent to

$$\mathbf{R}^* > 0 \Leftrightarrow \mathbf{W}_\mu^{-2} > \mathbf{W}_1^{-1} \Leftrightarrow \mathbf{W}_1 - \mathbf{W}_\mu \mathbf{W}_2 > \mu_1 \mathbf{I} + \mu_2 \mathbf{W}_2$$

$$\Leftrightarrow \lambda_1(\mathbf{W}_1 - \mu_2 \mathbf{W}_2) > \mu_1$$

(20)

where we used the standard tools of matrix analysis [31], [32]. When the IPC is redundant, $\mu_2 = 0$ from which (16) and

$$\mu_1^{-1} = m^{-1}(P_T + \text{tr}(\mathbf{W}_1^{-1}))$$

(21)

follow. It remains to establish (15). To this end, $\mu_2 = 0$ implies $\mu_1 > 0$ (active TPC) and $\lambda_1(\mathbf{W}_1) > \mu_1$ (from (20)) together with (21) implies 1st inequality in (15), 2nd inequality in (15) ensures that the IPC is redundant: $\mu_2 = 0$ and $\text{tr}(\mathbf{W}_2 \mathbf{R}^*) \leq P_T$. Indeed, the capacity under the joint constraints (TPC+IPC) cannot exceed that under the TPC alone, for which the optimal covariance $\mathbf{R}^*$ is as in (16). However, under 2nd inequality in (15), $\mathbf{R}^*$ satisfies $\text{tr}(\mathbf{W}_2 \mathbf{R}^*) \leq P_T$ and hence $\mathbf{R}^*$ is feasible under the joint constraints as well so that (i) both capacities are equal and (ii) $\mathbf{R}^*$ in (16) is also optimal under the joint constraints, which corresponds to $\mu_2 = 0$ (so that the IPC can be removed without affecting the capacity). (17) is obtained by using (16) in $C(\mathbf{R})$.

Next, we consider an interference-limited regime, where the TPC is redundant and hence the IPC is active.

**Proposition 4:** Let $\mathbf{W}_1, \mathbf{W}_2 > 0$ and $P_I$ be bounded as follows:

$$m \lambda_1(\mathbf{W}_2 \mathbf{W}_1^{-1}) - \text{tr}(\mathbf{W}_2 \mathbf{W}_1^{-1}) < P_I$$

$$\leq \frac{m}{\text{tr}(\mathbf{W}_2)}(P_I + \text{tr}(\mathbf{W}_1^{-1})) - \text{tr}(\mathbf{W}_2 \mathbf{W}_1^{-1})$$

(22)

then $\mu_1 = 0$, i.e. the TPC is redundant, $\mathbf{R}^*$ is of full-rank and is given by:

$$\mathbf{R}^* = \mu_2^{-1} \mathbf{W}_2^{-1} - \mathbf{W}_1^{-1}$$

(23)

where $\mu_2^{-1} = m^{-1}(P_I + \text{tr}(\mathbf{W}_2 \mathbf{W}_1^{-1}))$. The capacity can be expressed as

$$C = m \log \frac{P_I + \text{tr}(\mathbf{W}_2 \mathbf{W}_1^{-1})}{m} + \log |\mathbf{W}_1|$$

(24)

**Proof:** When the TPC is redundant, $\mu_1 = 0$ and (23) with

$$\mu_2^{-1} = m^{-1}(P_I + \text{tr}(\mathbf{W}_2 \mathbf{W}_1^{-1}))$$

(25)

follow from (10) in the same way as for Proposition 3. 1st condition in (22) follows from $\mathbf{R}^* > 0$, which, using (20), is equivalent to

$$\lambda_m(\mathbf{W}_2^{-1} \mathbf{W}_1) > \mu_2 = m(P_I + \text{tr}(\mathbf{W}_2 \mathbf{W}_1^{-1}))^{-1}$$

(26)

2nd condition in (22) ensures that the TPC is redundant since the Tx power is sufficiently large: $\mu_1 = 0$ and $\text{tr}(\mathbf{R}^*) \leq P_T$.

It is clear from (17) and (24) that the latter represents an interference-limited scenario while the former - a Tx power-limited one. In fact, (24) can be obtained from (17) via the substitutions $P_T \rightarrow P_I$ and $\mathbf{W}_1 \rightarrow \mathbf{W}_I = \mathbf{W}_2^{-1/2} \mathbf{W}_1 \mathbf{W}_2^{-1/2}$, where $\mathbf{W}_I$ plays a role of the equivalent channel and $\mathbf{W}_2^{-1/2}$ is a “whitening” filter at the transmitter. Note also that (17) and (24) coincide if $P_T = P_I$ and $\mathbf{W}_2 = \mathbf{I}$, as they should.

It follows from the proof of Proposition 3 that a general full-rank solution is given by

$$\mathbf{R}^* = \mathbf{W}_\mu^{-2} - \mathbf{W}_1^{-1}$$

(27)

and the respective capacity is $C = \log |\mathbf{W}_1^{-1} \mathbf{W}_1 \mathbf{W}_\mu^{-1}|$, which holds if $\lambda_m(\mathbf{W}_1 - \mu_2 \mathbf{W}_2) > \mu_1$. If neither (15) nor (22) hold, then both power constraints are active: $\mu_1, \mu_2 > 0$; closed-form solutions for $\mu_1, \mu_2$ in this case are not known and they have to be found numerically from (11). Note that $\mathbf{R}^*$ in (27) generalizes the respective WF full-rank solution $\mathbf{R}_W = \mu_1^{-1} \mathbf{I} - \mathbf{W}_1^{-1}$ (no IPC) to the interference-constrained environment, where $\mathbf{W}_\mu^{-2}$ takes the role of $\mu_1^{-1} \mathbf{I}$ and hence $\mathbf{W}_\mu^{-1}$ serves as a preceding “whitening” filter at the transmitter.

While the optimal covariance in (16) is the same as the standard WF over the eigenvectors of $\mathbf{W}_1$ (subject to the power constraint only), its range of validity is different: while the standard WF solution is of full rank under only the lower bound in (15) (this can be obtained by setting $P_I = \infty$, so that only the lower bound remains), i.e. the WF optimal covariance is of full rank for all sufficiently high power/SNR, the interference constraint also imposes the upper bound in (15) and hence its full-rank solution will not hold if $P_T$ is too high, a remarkable difference to the standard WF.

Next, we explore the case where $\mathbf{W}_2$ is of rank 1. This models the case when a primary user has a single-antenna receiver or when its channel is a keyhole channel, see e.g. [33], [34].
Proposition 5: Let \( W_1 \) be of full rank and \( W_2 \) be of rank-1, so that \( W_2 = \lambda_2 u_2 u_2^\dagger \), where \( \lambda_2 > 0 \) and \( u_2 \) are its active eigenvalue and eigenvector. If
\[
P_T \geq P_{I,th} = m^{-1} \lambda_2 (P_T + \text{tr}(W_1^{-1})) - \lambda_2 u_2^\dagger W_1^{-1} u_2
\]
\[
P_T > m \lambda_1 (W_1^{-1}) - \text{tr}(W_1^{-1})
\]
then the IPC is redundant, the optimal covariance is of full rank and is given by the standard WF solution,
\[
R^* = R_{WF}^* = \mu_1^{-1} W_1^{-1} - \alpha u_2 u_2^\dagger
\]
where \( \alpha = \mu_1^{-1} - (\mu_1 + \lambda_2 \mu_2)^{-1} \), and \( \mu_1, \mu_2 > 0 \) are found from (11):
\[
\mu_1 = (P_T - \lambda_2 P_T - u_2^\dagger W_1^{-1} u_2 + tr(W_1^{-1}))^{-1} (m - 1)
\]
\[
\mu_2 = (P_T + \lambda_2 u_2^\dagger W_1^{-1} u_2)^{-1} - \lambda_2^{-1} \mu_1
\]
Proof: See Appendix. □

Note that the 1st two terms in (32) represent the standard WF solution while the last term is a correction due to the IPC, which is reminiscent of a partial null forming in an adaptive antenna array, see e.g. [36]. Hence, partial null forming is also optimal from information-theoretic perspective in this case.

V. INACTIVE CONSTRAINTS

It is straightforward to see both constraints cannot be inactive at the same time (since the capacity is a strictly increasing function of the Tx power without the IPC). In this section, we explore the scenarios when one of the two constraints is redundant.2

Note that, unlike the standard WF where the TPC is always active, it can be inactive under the IPC, which is ultimately due to the interplay of interference and power constraints. The following proposition explores this in some details.

Proposition 6: The TPC is redundant only if \( N(W_2) \in N(W_1) \) and is active otherwise. In particular, it is active (for any \( P_T \) and \( P_I \)) if \( r(W_1) > r(W_2) \), e.g. if \( W_1 \) is full-rank and \( W_2 \) is rank-deficient.

Proof: Use (89) and note that this condition is necessary for the TPC to be redundant (since the KKT conditions are necessary for optimality and \( \mu_1 = 0 \) is also necessary for the TPC to be redundant). Now, if \( r(W_1) > r(W_2) \), then
\[
dim(N(W_2)) = m - r(W_2)
\]
\[
geq m - r(W_1) = \dim(N(W_1))
\]
where \( \dim(N) \) is the dimensionality of \( N \), and hence \( N(W_2) \in N(W_1) \) is impossible so that the TPC is active. □

If \( W_2 \) is full-rank, a more specific result can be established.

Proposition 7: If \( W_2 \) is redundant, then the optimal covariance \( R^* \) is as follows
\[
R^* = W_2^{-\frac{1}{2}} (\mu_2^{-1} W_2^{-\frac{1}{2}} W_1^{-1} W_2^{-\frac{1}{2}}) W_2^{-\frac{1}{2}},
\]
and \( \mu_2 > 0 \) is
\[
\mu_2^{-1} = \frac{1}{r_+} P_T + \frac{1}{r_+} \sum_{i=1}^{r_+} \lambda_{bi}^{-1}.
\]
where \( \lambda_{bi} = \lambda_i(W_2^{-1} W_1) \) and \( r_+ = r(R^*) \) is the rank of the optimal covariance matrix, which can be found as the largest solution of the following inequality:
\[
P_T > \sum_{i=1}^{r_+} (\lambda_{bi}^{-1} - \lambda_{bi}^{-1})
\]
In particular, this holds if
\[
P_T \geq \lambda_m^{-1}(W_2) P_T
\]
Proof: Let \( R^* \) be as in (35). Then, \( C \leq C(R^*) \), since \( C(R^*) \) is the capacity under the IPC only (this follows since (35) is a special case of (10) with \( \mu_1 = 0 \)). On the other hand, if \( P_T \geq tr(R^*) \), then \( R^* \) is feasible under the joint (IPC+TPC) constraints and hence \( C \geq C(R^*) \), which proves the equality \( C = C(R^*) \) and hence \( R^* \) is optimal. (36) follows from \( tr(W_2 R^*) = P_T \):
\[
P_T = tr(W_2 R^*) = tr(\mu_2^{-1} I - W_2^{-\frac{1}{2}} W_1^{-1} W_2^{-\frac{1}{2}})
\]
\[
= \sum_{i=1}^{r_+} (\mu_2^{-1} - \lambda_{bi}^{-1})
\]
(37) ensures, from (35), that the rank of \( R^* \) is \( r_+ \). Further note that (38) implies \( P_T \geq tr(R^*) \):
\[
P_T \geq \lambda_m^{-1}(W_2) P_T = \lambda_m^{-1}(W_2) tr(W_2 R^*) \geq tr(R^*)
\]
since \( W_2 \geq \lambda_m^{-1}(W_2) I \), but the converse is not true. □

Note that Proposition 7 gives a closed-form solution (including dual variables) for an optimal signaling strategy in the interference-limited regime, when the TPC is redundant and hence the IPC is active. The equivalent channel \( W_1 = W_2^{-\frac{1}{2}} W_1^{-1} W_2^{-\frac{1}{2}} \) plays a prominent role in this solution. The following proposition gives a sufficient and necessary condition for the IPC to be redundant (and hence the TPC is automatically active).

Proposition 8: The IPC is redundant and hence the standard WF solution is optimal, \( R^* = R_{WF}^* \), if and only if
\[
P_T \geq tr(W_2 R_{WF}^*)
\]
in which case the TPC is active: \( tr(R^*) = P_T \). In particular, this holds if
\[
P_T \geq \lambda_1(W_2) P_T
\]
Proof: Note that, under the joint (IPC+TPC) constraint, \( C \leq C(R_{WF}^*) \) (since \( C(R_{WF}^*) \) is attained by relaxing the

---

2“inactive” implies “redundant” but the converse is not true: for example, inactive TPC means \( tr(R^*) < P_T \) and this implies \( \mu_1 = 0 \) (from complementary slackness) so that it is also redundant (can be omitted without affecting the capacity), but \( \mu_1 = 0 \) does not imply \( tr(R^*) < P_T \) since \( tr(R^*) = P_T \) is also possible in some cases.
IPC and retaining the TPC only, which cannot decrease the optimum). On the other hand, under the stated conditions, \( R_{WF} \) is feasible under the joint constraints and hence the upper bound is achieved, \( C = C(R_{WF}) \) and \( R_{WF} \) is optimal. It is straightforward to see that (42) implies (41), since \( W_2 \leq \lambda_1 (W_2)I \) (note that (42) is sufficient but not necessary). □

While Proposition 8 gives sufficient and necessary conditions for the IPC to be redundant, they depend on \( P_T \) and \( P_I \). The next proposition gives a sufficient condition which is independent of \( P_T \) and \( P_I \) for the IPC to be redundant, and the standard WF solution is optimal: \( R^* = R_{WF} \).

**Proposition 9:** Let \( R(W_2) \in \mathcal{N}(W_1) \). Then, the IPC is redundant for any \( P_I \) and \( P_T \) and the standard WF solution is optimal: \( R^* = R_{WF} \).

**Proof:** As in Proposition 8, \( C \leq C(R_{WF}) \). It is straightforward to verify that any active eigenvector of \( R_{WF} \) is orthogonal to \( \mathcal{N}(W_1) \) and hence, due to \( R(W_2) \in \mathcal{N}(W_1) \), \( W_2 R_{WF} = 0 \), i.e., the IPC is redundant and \( R_{WF} \) is feasible (for any \( P_T \) and \( P_I \)) and hence \( C(R_{WF}) \leq C \), from which the desired result follows. □

It follows from this Proposition that ZF transmission is optimal in this case, for any \( P_I \) and \( P_T \), so that there is no loss of optimality in using this popular signalling technique.

**VI. RANK-1 SOLUTIONS**

In this section, we explore the case when \( W_1 \) is rank-one. As we show below, beamforming is optimal in this case. A practical appeal of this is due to its low-complexity implementation. Further, rank-one \( W_1 \) is also motivated by single-antenna mobile units while the base station is equipped with multiple antennas, or when the MIMO propagation channel is of degenerate nature resulting in a keyhole effect, see e.g., [33], [34].

We begin with the following result which bounds the rank of optimal covariance in any case.

**Proposition 10:** If the TPC is active or and \( W_2 \) is full-rank, then the rank of the optimal covariance \( R^* \) of the problem (P1) in (2) under the constraints in (8) is bounded as follows:

\[
r(R^*) \leq r(W_1) \tag{43}
\]

If the TPC is redundant and \( W_2 \) is rank-deficient, then there exists an optimal covariance \( R^* \) of (P1) under the constraints in (8) that also satisfies this inequality.\(^3\)

**Proof:** See Appendix. □

The following are immediate consequences of Proposition 10.

**Corollary 2:** If \( W_2 \) is of full-rank or and if the TPC is active, then the optimal covariance \( R^* \) is of full-rank only if \( W_1 \) is of full-rank (i.e. rank-deficient \( W_1 \) ensures that \( R^* \) is also rank-deficient).

**Corollary 3:** If \( r(W_1) = 1 \), then \( r(R^*) = 1 \), i.e. beamforming is optimal.

Note that this rank (beamforming) property mimics the respective property for the standard WF. However, while signalling on the (only) active eigenvector of \( W_1 \) is optimal under the standard WF (no IPC), it is not so when the IPC is active, as the following result shows. To this end, let \( W_1 = \lambda_1 u_1 u_1^\dagger \), i.e. it is rank-1 with \( \lambda_1 > 0 \), \( u_1 \) be the (only) active eigenvalue and eigenvector; \( \gamma_I = P_I/P_T \) be the “interference-to-signal” ratio, and

\[
\gamma_1 = \frac{u_1^\dagger W_2 u_1}{u_1^\dagger (W_2^*)^2 u_1}, \quad \gamma_2 = u_1^\dagger W_2 u_1 \tag{44}
\]

where \( W_2^\dagger \) is Moore-Penrose pseudo-inverse of \( W_2 \); \( W_2^2 = W_2^{-1} \) if \( W_2 \) is full-rank [31].

**Proposition 11:** Let \( W_1 \) be rank-1.

1. If \( \gamma_I < \gamma_1 \), then the TPC is redundant and the optimal covariance can be expressed as

\[
R^* = P_T W_2 W_2^\dagger W_2^\dagger u_1 W_2 W_2^\dagger u_1^\dagger \tag{45}
\]

The capacity is

\[
C = \log(1 + \lambda_1 \alpha P_T) \tag{46}
\]

where \( \alpha = \gamma_I u_1^\dagger W_2^\dagger u_1 \).\(^2\)

2. If \( \gamma_I \geq \gamma_2 \), then the IPC is redundant and the standard WF solution applies:

\[
R^* = P_T u_1 u_1^\dagger \tag{47}
\]

3. If \( \gamma_1 \leq \gamma_I < \gamma_2 \), then both constraints are active. The optimal covariance is

\[
R^* = P_T \frac{W_2 u_1 u_1^\dagger W_2^\dagger}{u_1^\dagger W_2^\dagger u_1} \tag{48}
\]

where \( W_2 = I + \mu_2 W_2 \), and \( \mu_2 > 0 \) is found from the IPC:

\[
tr(W_2 R^*) = P_T. \quad \text{The capacity is as in (46) with}
\]

\[\alpha = (u_1^\dagger W_2^\dagger u_1)^2 |W_2^\dagger u_1|^{-2} \leq 1 \tag{49}\]

with equality if and only if \( u_1 \) is an eigenvector of \( W_2 \).

**Proof:** See Appendix. □

Note that the optimal signalling in case 1 is along the direction of \( W_2^\dagger u_1 \) and not that of \( u_1 \) (unless \( u_1 \) is also an eigenvector of \( W_2 \)), as would be the case for the standard WF with redundant IPC. In fact, \( W_2^\dagger \) plays a role of a “whitening” filter here. Similar observation applies to case 3, with \( W_2 \) replaced by \( W_2 \). \( \alpha \) in Proposition 11 quantifies power loss due to enforcing the IPC; \( \alpha = 1 \) means no power loss.

**VII. IDENTICAL EIGENVECTORS**

In this section, we consider a scenario where \( W_1 \) and \( W_2 \) have the same eigenvectors. This may be the case when the scattering environment around the base station (the Tx) is the same as seen from the Rx and primary user \( U \). In this case, the general solution in Theorem 1 significantly simplifies to the following:

\[
R^* = U \Lambda^* U^+ \tag{49}
\]

where unitary matrix \( U \) collects eigenvectors of \( W_{1(2)} \) and diagonal matrix \( \Lambda^* = \text{diag}\{\lambda_i^*\} \) collects the eigenvalues of \( R^* \):

\[
\lambda_i^* = [(\mu_1 + \mu_2 \mu_1)^2 - \lambda_i]_+ \tag{50}
\]
where $\lambda_{1i(2i)}$ are the eigenvalues of $W_{1(2)}$. Dual variables $\mu_{1(2)} \geq 0$ are determined from the following:

$$
\mu_1 \left( \sum \lambda_i^* - P_T \right) = 0, \quad \mu_2 \left( \sum \lambda_{2i}^* - P_T \right) = 0 \quad (51)
$$

subject to $\sum_i \lambda_i^* \leq P_T$, $\sum_i \lambda_{2i}^* \leq P_T$. The capacity can be expressed as in (12) with $\lambda_{i1} = \lambda_{i1}(\mu_1 + \mu_2 \lambda_{2i})$.

Note that, in this case, signaling on the eigenmodes of the main channel $W_1$ is optimal, but power allocation is not given by the standard WF, unless the IPC is redundant ($\mu_2 = 0$). It follows from (50) that $\lambda_i^* = 0$ (power allocated to $i$-th eigenmode is zero) if $\mu_1 = 0$ (redundant TPC) and $\lambda_{2i} = 0$ (i-th eigenmode of 2nd channel $W_2$ is inactive), in addition to the standard WF property that $\lambda_i^* = 0$ if $\lambda_{i1} = 0$ under the active TPC ($\mu_1 > 0$).

In the context of massive MIMO systems under favorable propagation, see e.g. [37], $W_1$ and $W_2$ become diagonal matrices (and thus have the same eigenvectors), and so is $R^*$, i.e. independent signaling is optimal, and the solution in (50) gives the optimal power allocation in such setting. This significantly simplifies its implementation since numerical complexity of generic convex solvers can be prohibitively large for massive MIMO settings.

VIII. I TERATIVE BISECTION ALGORITHM

While Theorem 1 gives a closed-form solution for an optimal covariance $R^*$ up to dual variables, no closed-form solution is known for (11) in the general case; the sections above provided complete closed-form solutions in some special cases. In this section, we develop an iterative numerical algorithm to solve (11) in the general case in an efficient way and prove its convergence.

First, we consider the standard bisection algorithm [30]. Let $f(x)$ be a function with the following property: $f(x) \geq 0$ for any $x < x_0$ and $f(x) \leq 0$ for any $x > x_0$, where $x_0$ is a solution of $f(x) = 0$. Then, the following bisection algorithm (BA) can be used to solve $f(x) = 0$, where $x_t$, $x_u$ are the upper and lower bounds to $x_0$: $x_t \leq x \leq x_u$, and $\epsilon > 0$ is any desired accuracy. In fact, it is straightforward to show that this algorithm will converge in a finite number $N$ of steps such that

$$
N \leq \left\lceil \log_2 \left( \frac{x_u - x_t}{\epsilon} \right) \right\rceil \quad (52)
$$

where $[\cdot]$ denotes ceiling, so that the convergence is exponentially fast and hence the algorithm is very efficient [30].

**Algorithm 1 Bisection Algorithm (BA)**

- **Require:** $f(x)$, $x_t$, $x_u$, $\epsilon$
- **repeat**
  1. Set $x = \frac{1}{2}(x_t + x_u)$.
  2. If $f(x) < 0$, set $x_u = x$. Otherwise, set $x_t = x$.
  Terminate if $f(x) = 0$.
- **until** $|x_u - x_t| \leq \epsilon$.

An alternative stopping criteria for this algorithm is $|f(x)| \leq \epsilon$ and the two criteria are equivalent when $f(x)$ is continuous.

The BA can be used to solve for $\mu_1, \mu_2$ in (11) in an iterative way, as we show below. To this end, we need to establish lower and upper bounds to the solutions $\mu_1^*, \mu_2^*$ required by the BA.

**Proposition 12:** Let $\mu_1^*, \mu_2^*$ be solutions of (11), i.e. the optimal dual variables. They can be bounded as follows:

$$
0 \leq \mu_1^* \leq \mu_{1u} = m(P_T + \lambda_{11}^{-1}(W_1)^{-1})^{-1} \quad (53)
$$

$$
0 \leq \mu_2^* \leq \mu_{2u} = (P_T/r_2 + \lambda_m(W_2)/\lambda_1(W_1))^{-1} \quad (54)
$$

where $r_2$ is the rank of $W_2$ and $m$ is the number of Tx antennas.

**Proof:** From the KKT conditions in (81),

$$
(I + W_1 R)^{-1} W_1 R = \mu_1 R + \mu_2 W_2 R,
\quad \mu_1 P_T + \mu_2 P_T = tr((I + W_1 R)^{-1} W_1 R) \quad (55)
$$

Let $A$ be a matrix with positive eigenvalues, $\lambda_i(A) \geq 0$. Since

$$
\lambda_i((I + A)^{-1} A) = \lambda_i(A)(1 + \lambda_i(A))^{-1} \leq \lambda_i(A)(1 + \lambda_1(A))^{-1} \quad (56)
$$

it follows that

$$
tr((I + A)^{-1} A) \leq m\lambda_1(A)(1 + \lambda_1(A))^{-1} \quad (57)
$$

Now use $A = W_1 R$ to obtain

$$
tr((I + W_1 R)^{-1} W_1 R) \leq m\lambda_1(W_1 R)(1 + \lambda_1(W_1 R))^{-1} \leq mP_T((\lambda_{11}^{-1} + P_T)^{-1} \quad (58)
$$

where $\lambda_{11} = \lambda_1(W_1)$ and $\lambda_1(W_2) \leq \lambda_{11} P_T$, so that 2nd inequality in (53) follows from (55). Let $\lambda_{2m} = \lambda_m(W_2)$.

Using (10) under active IPC,

$$
P_t = tr(W_2 R^*) \leq tr(W_2(W_2^*)^2((1 + \epsilon) - 1 - \lambda_{11}^{-1}(1 + \epsilon) \lambda_{2m}))$$

$$
\leq r_2(\mu_1^2 - \lambda_{2m} \lambda_{11}^{-1}) \quad (59)
$$

from which 2nd inequality in (54) follows, where we have used

$$
tr(W_2(W_2^*)^2) \leq \mu_2^{-1} tr(W_2 W_2^*) = \mu_2^{-1} r_2 \quad (60)
$$

and $W \geq \lambda_m(W)I$ for any Hermitian $W$. If the IPC is inactive, $\mu_2 = 0$ and the inequality holds in obvious way. □

To proceed further, let

$$
x_{\epsilon} = BA[f(x), x_t, x_u, \epsilon] \quad (61)
$$

formally denote an $\epsilon$-accurate solution of $f(x) = 0$ given by the BA and let

$$
f_1(\mu_1, \mu_2) = \mu_1 \left(tr(R^*(\mu_1, \mu_2)) - P_T \right)
$$

$$
f_2(\mu_1, \mu_2) = \mu_2 \left(tr(W_2 R^*(\mu_1, \mu_2)) - P_T \right) \quad (62)
$$

where $R^*(\mu_1, \mu_2)$ denotes $R^*$ in (10) for given $\mu_1, \mu_2$. Then, the optimal dual variables $\mu_1^*, \mu_2^*$ satisfy $f_1(\mu_1^*, \mu_2^*) = 0$ and $f_2(\mu_1^*, \mu_2^*) = 0$. For a given $\mu_2^*$, one could use the BA to formally express $\mu_1^*$ as

$$
\mu_1^* = BA[f(x), f_1(x, \mu_2^*), \mu_1, \mu_{1u}, 0] \quad (63)
$$

where, from (53), $\mu_1 = 0$, and likewise for $\mu_2^*$ (since the convergence of the BA is exponentially fast, the accuracy $\epsilon$ is not relevant).
can be set to be arbitrary small in practice so that we disregard here this small inaccuracy by setting $\epsilon = 0$ to simplify the analysis; numerical experiments support this approach. The following proposition shows that $f_1(x, \mu_2)$, $f_2(\mu_1, x)$ have the property needed for the convergence of the BA as stated above. To this end, let $P_2(\mu_1, \mu_2) = tr(R^*(\mu_1, \mu_2))$, $P_2(\mu_1, \mu_2) = tr(W_2 R^*(\mu_1, \mu_2))$, i.e. the transmit and interference powers for given $\mu_1, \mu_2$.

**Proposition 13:** Let $\mu_1$ be a solution of $f_1(x, \mu_2) = 0$ for a given $\mu_2$ and subject to $P_1(x, \mu_2) \leq P_2$. Then, $f_1(\mu, \mu_2) \geq 0$ for any $\mu < \mu_0$ and $f_1(\mu_0, \mu_2) \leq 0$ for any $\mu > \mu_0$. Likewise, if $\mu_2$ is a solution of $f_2(\mu_1, x) = 0$ for a given $\mu_1$ and subject to $P_2(\mu_1, x) \leq P_1$, then $f_2(\mu_1, \mu_2) \geq 0$ for any $\mu_2 < \mu_0$ and $f_2(\mu_1, \mu_2) \leq 0$ for any $\mu_2 > \mu_0$.

**Proof:** See the extended version of this paper [38].

Thus, this proposition shows that the BA can be used to solve $f_1(x, \mu_2) = 0$ for a given $\mu_2$ and likewise for $f_2(\mu_1, x) = 0$. Unfortunately, neither of the optimal dual variables is known in advance. Hence, we propose the following iterative bisection algorithm (IBA) which finds optimal dual variables without such advance knowledge.

**Algorithm 2 Iterative Bisection Algorithm (IBA)**

**Require:** $f_1(\mu_1, \mu_2)$, $f_2(\mu_1, \mu_2)$, $\mu_{1u}$, $\mu_{2u}$, $\delta$

1. Set $\mu_{20} = 0$, $k = 1$.

   repeat

   2. Set $\mu_{1k} = BA[f_1(x, \mu_{2(k-1)}), 0, \mu_{1u}, \delta]$.

   3. Set $\mu_{2k} = BA[f_2(\mu_{1k}, x), 0, \mu_{2u}, \delta]$.

   4. $k := k + 1$.

   until stopping criterion is met.

Note that the BA used in steps 2 and 3 will converge, as follows from Proposition 13. A possible stopping criteria for this algorithm is $|f_1(\mu_{1k}, \mu_{2k})| \leq \epsilon$ or when a number of steps exceeds maximum $k_{max}$. The following proposition shows that the IBA generates converging sequences of dual variables $\{\mu_{1k}\}$, $\{\mu_{2k}\}$ under a mild technical condition.

**Proposition 14:** The sequences $\{\mu_{1k}\}_{k=1}^{\infty}$, $\{\mu_{2k}\}_{k=1}^{\infty}$ generated by the IBA above converge if $\delta = 0$ and $P_1(\mu_1, \mu_2) < P_2(\mu_1, \mu_2)$ for any $\mu_2 < \mu_0$ and $f_2(\mu_1, \mu_2) \leq 0$ for any $\mu_2 > \mu_0$.

**Proof:** See the extended version of this paper [38].

The following proposition shows that any stationary (and hence convergence) point of the IBA solves the dual optimality conditions in (11).

**Proposition 15:** Any stationary point of the IBA is a solution of (11) if $\delta = 0$. Hence, the IBA converges to a solution of (11) under the conditions of Proposition 14.

**Proof:** Let $\mu_{1s}$, $\mu_{2s}$ be a stationary point of the IBA, so that

\[
\begin{align*}
\mu_{1s} &= BA[f_1(x, \mu_{2s}), 0, \mu_{1u}, 0], \\
\mu_{2s} &= BA[f_2(\mu_{1s}, x), 0, \mu_{2u}, 0]
\end{align*}
\]

(64)

It follows from 1st equality that $f_1(\mu_{1s}, \mu_{2s}) = 0$ and $f_2(\mu_{1s}, \mu_{2s}) = 0$ from 2nd one. Thus, $\mu_{1s}$, $\mu_{2s}$ solves (11). Since a convergence point is stationary, it follows that the IBA converges to a solution of (11).

While the analytical convergence results above are limited to $\delta = 0$, $\delta > 0$ is used in practice. Since the BA converges exponentially fast, very small $\delta$ can be selected in the IBA without significant increase in computational complexity of each step and hence the analysis serves as a reasonable approximation (due to the continuity of the problem and functions involved). Furthermore, numerous numerical experiments indicate that the IBA always converges, even when the conditions 1-3 of Proposition 14 are not met (we were not able to observe a single case where it did not). In the majority of the studied cases, a small to moderate number of IBA iterations (1...50) is needed to achieve a high accuracy of $10^{-5}$, while up to 250 iterations are required in some exceptional cases with $\epsilon = 10^{-10}$ (which is hardly required in practice).

**IX. EXTENSION TO MULTI-USER ENVIRONMENTS**

In a typical multi-user environment, there are multiple users to which interference has to be limited, so that the problem in (2) is solved under the following constraint set:

\[
S_R = \{R : R \succeq 0, tr(R) \leq P_T, tr(W_2 R) \leq P_{1k}, k = 1..K\},
\]

(65)

where $W_{2k} = H_{2k}^+ H_{2k}$ and $P_{1k}$ represent channel to user $k$ and respective interference constraint power, $K$ is the number of users, see Fig. 2. Using the same approach as in Theorem 1, it is straightforward to see that Theorem 1 applies with

\[
W_\mu = (\mu_1 I + \sum_k \mu_{2k} W_{2k})^{\frac{1}{2}}
\]

(66)

where dual variables are found from the following system of non-linear equations:

\[
\begin{align*}
\mu_1 (tr(R^*) - P_T) &= 0, \\
\mu_{2k} (tr(W_{2k} R^*) - P_{1k}) &= 0, \quad k = 1..K
\end{align*}
\]

(67)

subject to $\mu_1, \mu_{2k} \geq 0$, $tr(R^*) \leq P_T$, $tr(W_{2k} R^*) \leq P_{1k}$. In particular, the iterative bisection algorithm of Section VIII can be used with a proper extension to accommodate multiple users.

In this multi-user setting, Proposition 1, Corollary 1 and Proposition 6 hold with the substitution $W_2 \rightarrow \sum_k W_{2k}$, i.e.
to solve (11) with the accuracy \( \epsilon = 10^{-5} \) vs. \( P_T \); \( W_1 \) and \( W_2 \) are as in (73), \( P_I = 1 \).

(i) the capacity grows unbounded if and only if

\[
\mathcal{N}\left( \sum_k W_{2k} \right) \notin \mathcal{N}(W_1) \tag{68}
\]

and (ii) the TPC is redundant only if

\[
\mathcal{N}\left( \sum_k W_{2k} \right) \in \mathcal{N}(W_1) \tag{69}
\]

and is active otherwise. In particular, it is active (for any \( P_T \) and \( P_{Ik} \)) if \( r(W_1) > r(\sum_k W_{2k}) \), e.g. if \( W_1 \) is full-rank and \( \sum_k W_{2k} \) is rank-deficient.

Proposition 2 holds with the substitution \( W_2 \to \sum_{k \in K_0} W_{2k} \), where \( K_0 = \{ k : P_{Ik} = 0 \} \) is a set of all primary users requiring no interference (\( P_{Ik} = 0 \)), so that the capacity is zero if and only if \( P_{Ik} = 0 \) for some \( k \) and

\[
\mathcal{N}\left( \sum_{k \in K_0} W_{2k} \right) \in \mathcal{N}(W_1). \tag{70}
\]

One may also consider the total (rather than individual) interference power constraint so that

\[
S_R = \{ R : R \geq 0, \ tr(R) \leq P_T, \sum_k tr(W_{2k}R) \leq P_I \} \tag{71}
\]

In this case, all the results apply with the substitution

\[
W_2 \to \sum_k W_{2k} \tag{72}
\]

X. EXAMPLES

In this section, we present some numerical results that illustrate the analytical results above as well as the performance of the IBA.

**Example 1:** In this example, \( P_I = 1 \) and

\[
W_1 = \begin{bmatrix} 1 & 0 \\ 0 & 0.5 \end{bmatrix}
\]

\[
W_2 = \begin{bmatrix} 1 & -0.5 \\ -0.5 & 1 \end{bmatrix} \quad \text{or} \quad \begin{bmatrix} 1 & 0.5j \\ -0.5j & 1 \end{bmatrix} \tag{73}
\]

Fig. 3 shows the number of iterations of the IBA required to solve (11) with the accuracy \( \epsilon = 10^{-5} \) vs. \( P_T \); the optimal dual variables \( \mu_1^*, \mu_2^* \) as well as the actual Tx and interference powers \( P_1 = tr(R^*(\mu_1^*, \mu_2^*)) \) and \( P_2 = tr(W_2R^*(\mu_1^*, \mu_2^*)) \), respectively, are also shown; the same results hold for both \( W_2 \) in (73).

Fig. 3 shows the number of iterations of the IBA required to solve (11) with the accuracy \( \epsilon = 10^{-5} \) vs. \( P_T \); the optimal dual variables \( \mu_1^*, \mu_2^* \) as well as the actual Tx and interference powers \( P_1 = tr(R^*(\mu_1^*, \mu_2^*)) \) and \( P_2 = tr(W_2R^*(\mu_1^*, \mu_2^*)) \), respectively, are also shown; the same results hold for both \( W_2 \) in (73). Note the transition from the Tx power-limited regime (inactive IPC) to the interference-limited regime (inactive TPC) as \( P_T \) increases, which is visible when the respective dual variable sharply decreases to 0. In particular, the IPC is inactive when \( P_T < 1 \) and the TPC is inactive when \( P_T > 1.9 \), while both constraints are active otherwise. As \( P_T \) increases, the IPC becomes active at about \( P_T \approx 1 \), at which point the required number of iteration sharply increases from 1 to 36 and then to 82, gradually decreasing to a small number of 2. When the IPC is inactive, the number of iterations is 1, in agreement with Proposition 14. As this example demonstrates, anyone of the constraints can be inactive depending on the \( P_T, P_I \) and channel matrices.

Fig. 4 shows the capacity under the joint (TPC+IPC) constraints for the channel of Fig. 3, along with the capacities under the TPC \( (C_{WF}) \), IPC \( (C_{IPC}) \) and joint TPC+IPC \( (C) \) alone. Note that \( C \) is upper bounded in general by \( C_{WF} \) and \( C_{IPC} \),

\[
C \leq \min\{C_{WF}, C_{IPC}\} \tag{74}
\]

and that this bound is tight: if the IPC is inactive (power-limited regime), then \( C = C_{WF} \), and if the TPC is inactive (interference-limited regime), then \( C = C_{IPC} \), so that the inequality is strict only in a (small) transition region (when both constraints are active simultaneously) and hence the following approximation can be used over the entire range of \( P_T \):

\[
C \approx \min\{C_{WF}, C_{IPC}\} \tag{75}
\]

Note also that the capacity does not grow unbounded, in agreement with Proposition 1, since \( W_2 \) is full-rank, so that \( \mathcal{N}(W_2) \) is empty and hence \( \mathcal{N}(W_2) \notin \mathcal{N}(W_1) \). This changes significantly if \( W_2 \) is rank-deficient: the TPC is always active and the capacity grows unbounded, as the next example shows.
Example 2: In this example

\[
W_2 = \begin{bmatrix} 1 & -1 \\ -1 & 1 \end{bmatrix} \quad \text{or} \quad \begin{bmatrix} 1 & j \\ j & 1 \end{bmatrix}
\]  

so that \(W_2\) is rank-deficient, and \(W_1\) is as in Example 1. As Fig. 5 shows, the capacity grows unbounded for any \(P_1\), even small one, in agreement Corollary 1, since \(r(W_1) > r(W_2)\), or Proposition 1, since \(\mathcal{N}(W_2) \notin \mathcal{N}(W_1)\). In this case, the bound in (74) is tight and the approximation in (75) is valid at low SNR only, since \(C_{IPC} = \infty\) as \(W_1\) is full-rank while \(W_2\) is rank-deficient.

Comparing Fig. 5 to Fig. 4, one observes that while decreasing \(P_1\) decreases the capacity in both cases, the behaviour is qualitatively different: the capacity saturates and variations in \(P_1\) have major impact on the saturation level in Fig. 4 while the capacity grows unbounded in Fig. 5 and variations in \(P_1\) have moderate or small impact on its value (negligible if \(P_1 \leq 0.1\)). In both cases, the channel matrix of primary user has a major impact on the capacity at high power/SNR regime, while being negligible at low SNR. Its null space (or rank) determines the qualitative behaviour of the capacity at high power/SNR regime.

It should also be noted that the optimal covariance \(R^*\) is not diagonal, even though \(W_1\) is, when the IPC is active - a sharp distinction to the TPC constraint only, where \(R^*\) and \(W_1\) have the same eigenvectors so that diagonal \(W_1\) implies diagonal \(R^*\). Hence, introducing the IPC makes independent signaling sub-optimal for independent channels in general (unless \(W_2\) is also diagonal or if the IPC is redundant).

Example 3: Figure 6 shows typical performance for a larger number of antennas \(m = n_1 = n_2 = 10\) and randomly-generated \(H_1, H_2\) with i.i.d. complex Gaussian entries (zero mean and unit variance). Comparing it to Fig. 3, it is clear that the general tendencies are the same, while numerical values are somewhat different: the transition region from the TPC-dominated regime to the IPC-dominated one is somewhat larger \((0.048 < P_T < 1.9)\), where both constraints are active simultaneously), while the maximum number of iterations is actually smaller compared to Fig. 3. It takes only 1 iteration to converge in the TPC-dominated regime (small \(P_T\)) and only 2 iterations in the IPC-dominated one (large \(P_T\)); the number of iterations is moderately large in the transition region. It should be noted that this performance is typical for many random channel realizations, with only minor variations from realization to realization when \(m, n\) are large.

Example 4: To demonstrate that \(R^*\) is not necessarily unique under the IPC (a stark difference to the standard WF solution, where the optimal covariance is always unique, unless the capacity is zero), let us consider the following channel:

\[
W_1 = \text{diag}\{w_1, 0\}, \quad W_2 = \text{diag}\{w_2, 0\}, \quad w_2P_T > P_I
\]  

It is straightforward to see that an optimal covariance is

\[
R^* = \text{diag}\{w_2^{-1}P_I, a\}, \quad 0 \leq a \leq P_T - w_2^{-1}P_I
\]  

so that it is not unique, but the capacity is:

\[
C = \ln(1 + w_1w_2^{-1}P_I)
\]

for any \(a\). While \(r(R^*) = 1\) for \(a = 0\) so that (43) does hold, this is not the case for \(a > 0\), which is in stark contrast to the standard WF solution, where (43) always holds. Note that this unusual property disappears if \(w_2P_T \leq P_I\), in which case the IPC is redundant and the standard WF solution applies:

\[
R^* = \text{diag}\{P_T, 0\}, \quad C = \ln(1 + w_1P_T).
\]

It should be noted that all the above numerical examples are representative in the sense that similar tendencies also hold for different numbers of antennas and different channels.

XI. CONCLUSION

Optimal signaling over the Gaussian MIMO channel is considered under the total TX and interference power constraints. The closed-form solution for an optimal Tx covariance is presented in the general case (up to dual variables). The iterative bisection algorithm is developed to evaluate numerically the
dual variables in the general case and its convergence is proved for some special cases. A number of more explicit closed-form solutions (with optimal dual variables) are obtained, including full-rank, rank-1 (beamforming) and when the channels have the same eigenvectors. Sufficient and necessary conditions for the TPC and IPC being active/inactive (or redundant) are given and their interplay is investigated. It is pointed out that the TPC and IPC can be active simultaneously so that neither condition can be absorbed into another in general, as was sometimes suggested in the literature. Null spaces of the channel matrices of primary users have a major impact on the capacity at high SNR while being negligible at low SNR. These analytical results can serve as building blocks for the analysis, design and optimization of multi-user MIMO networks in interference-limited environments, as in e.g. 5G scenarios with aggressive frequency re-use, HetNets and licensed/licenused usage to improve spectral efficiency.

APPENDIX

A. Proof of Theorem 1

Since the problem is convex and Slater’s condition holds, the KKT conditions are both sufficient and necessary for optimality [30]. They take the following form:

\[-(I + W_1 R)^{-1} W_1 - M + \mu_1 I + \mu_2 W_2 = 0\]  \hspace{1cm} (81)

\[MR = 0, \quad \mu_1 (tr(R) - P_T) = 0,\]  \hspace{1cm} (82)

\[\mu_2 (tr(W_2 R) - P_T) = 0,\]  \hspace{1cm} (83)

\[tr(R) \leq P_T, \quad tr(W_2 R) \leq P_I, \quad R \succeq 0\]  \hspace{1cm} (84)

where \(M\) is Lagrange multiplier responsible for the positive semi-definite constraint \(R \succeq 0\). We consider first the case of full-rank \(W_\mu\) (i.e. either \(\mu_1 > 0\) or/and \(W_2 > 0\)), so that \(W_\mu^{-1} = W^{-1}_\mu\). Let us introduce new variables: \(\tilde{R} = W_\mu R W_\mu\), \(W_1 = W_\mu^{-1} W_1 W_\mu^{-1}\), \(M = W_\mu^{-1} M W_\mu^{-1}\). It follows that \(\tilde{M} \tilde{R} = 0\) and (81) can be transformed to

\[(I + \tilde{W}_1 \tilde{R})^{-1} \tilde{W}_1 + \tilde{M} = I\]  \hspace{1cm} (85)

for which the solution is

\[\tilde{R} = (I - \tilde{M})^{-1} - \tilde{W}_1^{-1} = (I - \tilde{W}_1^{-1})_+\]  \hspace{1cm} (86)

(this can be established in the same way as for the standard WF in (4)). Transforming back to the original variables results in (10). (11) are complementary slackness conditions in (82); (12) follows, after some manipulations, by using \(R^*\) of (10) in \(C(R)\).

The case of singular \(W_\mu\) is more involved. It implies \(\mu_1 = 0\) so that \(W_\mu = \mu_2 W_2\). It follows from the KKT condition in (81) that, for the redundant TPC (\(\mu_1 = 0\)),

\[Q_1 (I + Q_1 R Q_1)^{-1} Q_1 + M = \mu_2 W_2\]  \hspace{1cm} (87)

where \(Q_1 = W_1^{1/2}\). Let \(x \in \mathcal{N}(W_2)\), i.e. \(W_2 x = 0\), then

\[x^+ Q_1 (I + Q_1 R Q_1)^{-1} Q_1 x + x^+ M x = 0\]  \hspace{1cm} (88)

so that \(x^+ M x = 0\) and \(Q_1 x = 0\), since \(M \geq 0\) and \(I + Q_1 R Q_1 > 0\). Thus, \(\mathcal{N}(W_2) \in \mathcal{N}(Q_1) = \mathcal{N}(W_1)\) and \(\mathcal{N}(W_2) \in \mathcal{N}(M)\), i.e.

\[\mathcal{N}(W_2) \in \mathcal{N}(W_1) \cap \mathcal{N}(M)\]  \hspace{1cm} (89)

and this condition is also necessary for the TPC to be redundant. Using (87), (89) and introducing new variables

\[\Lambda_2 = U_2^+ W_2 U_2, \quad \tilde{R} = U_2^+ R U_2,\]  \hspace{1cm} (90)

\[Q_1 = U_2^+ Q_1 U_2, \quad M = U_2^+ M U_2,\]  \hspace{1cm} (91)

where \(U_2\) is a unitary matrix of eigenvectors of \(W_2\), one obtains

\[\Lambda_2 = \begin{pmatrix} \Lambda_{2+} & 0 \\ 0 & 0 \end{pmatrix}, \quad \tilde{Q}_1 = \begin{pmatrix} Q_{1+} & 0 \\ 0 & 0 \end{pmatrix}\]  \hspace{1cm} (92)

\[\tilde{M} = \begin{pmatrix} M_+ & 0 \\ 0 & 0 \end{pmatrix}, \quad \tilde{R} = \begin{pmatrix} R_+ & R_{12} \\ R_{21} & R_{22} \end{pmatrix}\]  \hspace{1cm} (93)

where \(\Lambda_{2+} > 0\) is a diagonal matrix of strictly positive eigenvalues of \(W_2\), so that (87) can be transformed to

\[Q_1+(I + Q_1+ R_+ Q_{1+})^{-1} Q_1+ + M_+ = \mu_2 \Lambda_{2+} > 0\]  \hspace{1cm} (94)

Using

\[Q_1+(I + Q_1+ R_+ Q_{1+})^{-1} Q_1+ = (I + W_1+ R_+)^{-1} W_1+\]  \hspace{1cm} (95)

where \(W_1+ = Q_{1+}^2\) and adopting (89), (10), one obtains

\[R_+ = \Lambda_{2+}^{-\frac{1}{2}} (\mu_2^{-1} I - \Lambda_{2+}^{-\frac{1}{2}} W_1^{-1} \Lambda_{2+}^{-\frac{1}{2}}) + \Lambda_{2+}^{-\frac{1}{2}}\]  \hspace{1cm} (96)

Since

\[P_T \geq tr(R) = tr(\tilde{R}) \geq tr(R_+),\]

\[P_I \geq tr(W_2 R) = tr(\Lambda_2 R_2) = tr(\Lambda_2 R_+),\]  \hspace{1cm} (97)

one can set, without loss of optimality, \(R_{22} = 0\), \(R_{12} = 0\), \(R_{21} = 0\), and transform (93) to

\[\tilde{R} = (\Lambda_2^T \Lambda_{2+}^{-\frac{1}{2}} (\mu_2^{-1} I - \Lambda_{2+}^{-\frac{1}{2}} W_1^{-1} \Lambda_{2+}^{-\frac{1}{2}}) + \Lambda_{2+}^{-\frac{1}{2}}\]  \hspace{1cm} (98)

and hence, as desired,

\[R = U_2 \tilde{R} U_2^T = W_\mu^T (I - W_\mu W_1^{-1} W_\mu) + W_\mu^T\]  \hspace{1cm} (99)

B. Proof of Proposition 1

To prove the “if” part, observe that \(\mathcal{N}(W_2) \notin \mathcal{N}(W_1)\) implies \(\exists u : W_2 u = 0, W_1 u \neq 0\). Now set \(R = P_T u u^+\), for which \(tr(R) = P_T, tr(W_2 R) = 0\), so it is feasible for any \(P_T, P_I\). Furthermore,

\[C \geq C(R) = \log(1 + P_T u^+ W_1 u) \to \infty\]  \hspace{1cm} (100)

as \(P_T \to \infty\), since \(u^+ W_1 u > 0\).

To prove the “only if” part, assume that \(\mathcal{N}(W_2) \in \mathcal{N}(W_1)\). This implies that \(\mathcal{R}(W_1) \in \mathcal{R}(W_2)\) (since \(\mathcal{R}(W)\) is the complement of \(\mathcal{N}(W)\) for Hermitian \(W\)). Let

\[W_k = U_{k+} \Lambda_k U_{k+}^+, \quad k = 1, 2\]  \hspace{1cm} (101)
where $U_{k+}$ is a semi-unitary matrix of active eigenvectors of $W_k$ and diagonal matrix $\Lambda_k$ collects its strictly-positive eigenvalues. Notice that, from the IPC,

$$
P_T \geq \text{tr}(W_2 R) = \text{tr}(\Lambda_2 U_{2+}^* R U_{2+})
\geq \lambda_{r_2} \text{tr}(U_{2+}^* R U_{2+})
$$

(99)

where $\lambda_{r_2} > 0$ is the smallest positive eigenvalue of $W_2$, so that

$$
\lambda_1(U_{2+}^* R U_{2+}) \leq P_1/\lambda_{r_2} < \infty
$$

(100)

for any $P_T$. On the other hand, $R(W_1) \in R(W_2)$ implies $\text{span}\{U_{1+}\} \in \text{span}\{U_{2+}\}$ and hence

$$
\lambda_1(U_{1+}^* R U_{1+}) \leq \lambda_1(U_{2+}^* R U_{2+}) \leq P_1/\lambda_{r_2} < \infty
$$

(101)

so that

$$
C(P_T) = \log |I + \Lambda_1 U_{1+}^* R^* U_{1+}| = \sum_i \log(1 + \lambda_i(A_1 U_{1+}^* R^* U_{1+}))
\leq m \log(1 + \lambda_1(W_1) \lambda_1(U_{1+}^* R^* U_{1+}))
\leq m \log(1 + \lambda_1(W_1) P_1/\lambda_{r_2}) < \infty
$$

(102)

for any $P_T$, as required.

C. Proof of Proposition 5

Start with the matrix inversion Lemma to obtain

$$
(\mu_1 I + \mu_2 \lambda_2 u_2 u_2^*)^{-1} = ((\mu_1 + \mu_2 \lambda_2)^{-1} - \mu_2^{-1} u_2 u_2^*) + \mu_1^{-1} I
$$

(103)

so that (32) follows from (27). Since $W_1$ is full-rank and $W_2$ is rank-1, it follows that the TPC is always active, $\mu_1 > 0$ and $\text{tr}(R) = \lambda_2$, from which one obtains

$$
m \mu_1^{-1} - \alpha - \text{tr}(W_1^{-1}) = P_T
$$

(104)

When the IPC is active, $\text{tr}(W_2 R) = P_T$, it follows that

$$
\lambda_2(\mu_1 + \mu_2 \lambda_2)^{-1} = \mu_1 + \mu_2 u_2^* W_1^{-1} u_2
$$

(105)

Solving (104) and (105) for $\mu_1$, one obtains 1st equality in (33); using it in (105) results in 2nd equality in (33). (31) and 1st inequality in (30) ensure that $R^* > 0$, since

$$
\mu_1^{-1} > \lambda_1(W_1^{-1}) + \alpha \geq \lambda_1(W_1^{-1} + \alpha u_2 u_2^*)
$$

(106)

where 1st inequality is due to 1st inequality in (30) and (105) while 2nd inequality is from $\lambda_1(A + B) \leq \lambda_1(A) + \lambda_1(B)$ where $A$, $B$ are Hermitian matrices (see e.g. [31]). It follows from (106) that $\mu_1^{-1} I > W_1^{-1} + \alpha u_2 u_2^*$ and hence $R^* > 0$, and that $\mu_1 > 0$, as required. 2nd inequality in (30) ensures that the IPC is active, $\mu_2 > 0$.

To obtain (29), observe that $R_{WF}$ is feasible under (28):

$$
\text{tr}(R_{WF}) = P_T, \text{tr}(W_2 R_{WF}) \leq P_1, R_{WF} > 0.
$$

(107)

Since it is a solution without the IPC (as the standard full-rank WF solution), it is also optimal under the IPC.

D. Proof of Proposition 10

We consider first the case when $W_\mu$ is full-rank, i.e. when either the TPC is active, $\mu_1 > 0$, or/and $W_2 > 0$. It follows from (81) that

$$
(I + W_1 R^*)^{-1} W_1 R^* = W_\mu R^*
$$

(108)

so that, since $(I + W_1 R)$ and $W_\mu$ are full-rank,

$$
r(R^*) = r(W_\mu R^*) = r(W_1 R^*)
\leq \min(r(W_1), \ r(R^*)) \leq r(W_1)
$$

(109)

The case of rank-deficient $W_\mu$ (i.e. when $\mu_1 = 0$ and $W_2$ is rank-deficient) is more involved. In this case, it follows from Proposition 6 that $N(W_2) = N(W_1)$ and hence $R(W_1) \in R(W_2)$ (if $W$ is Hermitian, $R(W)$ is the complement of $N(W)$), from which the following equivalency can be established, which is instrumental in the proof.

Proposition 16: If $W_2$ is rank-deficient and the TPC is redundant for the problem (P1) in (2) under the constraint in (8), then (P1) has the same value as the following problem (P2):

$$(P2) : \ max \tilde{C}(\tilde{R}) \ s.t. \ tr(\tilde{A}_2 \tilde{R}) \leq P_1, \ tr(\tilde{R}) \leq P_T$$

(110)

where $\tilde{C}(\tilde{R}) = |I + \tilde{W}_1 \tilde{R}|$, $\tilde{W}_1 = U_{2+}^* W_1 U_{2+}$, $\tilde{A}_2 = U_{2+}^* W_2 U_{2+} > 0$ is a diagonal matrix of strictly-positive eigenvalues of $W_2$ and $U_{2+}$ is a semi-unitary matrix whose columns are the corresponding active eigenvectors of $W_2$. Furthermore, an optimal covariance $R^*$ of (P1) can be expressed as follows:

$$
R^* = U_{2+} \tilde{R} U_{2+}^*
$$

(111)

where $\tilde{R}^*$ is a solution of (110):

$$
\tilde{R}^* = \tilde{A}_2^{-\frac{1}{2}} (\mu_2^{-1} I - \tilde{A}_2^{-1} \tilde{W}_1^{-1} \tilde{A}_2^{-\frac{1}{2}} + \tilde{A}_2^{-\frac{1}{2}})
$$

and $\mu_2 > 0$ is found from the IPC:

$$
\text{tr}(\mu_2^{-1} I - \tilde{A}_2^{-1} \tilde{W}_1^{-1} \tilde{A}_2^{-\frac{1}{2}} + \tilde{A}_2^{-\frac{1}{2}}) = P_T
$$

(112)

Proof: Let $R^*$ and $\tilde{R}^*$ be the solutions of (P1) and (P2) under the stated conditions and let $P_2 = U_{2+} U_{2+}^*$ be a projection matrix on the space spanned by the active eigenvectors of $W_2$, i.e. on $R(W_2)$. Note that $P_2 W_k P_2 = W_k$, $k = 1, 2$, since $R(W_1) \in R(W_2)$ under the stated conditions. Define $\tilde{R}' = U_{2+} R' U_{2+}$ and observe that

$$
P_T \geq \text{tr}(R^*) \geq \text{tr}(\tilde{R}')
$$

(114)

so that $\tilde{R}'$ is feasible for (P2) and hence

$$
\tilde{C}(\tilde{R}') \geq \tilde{C}(\tilde{R}') = \log |I + \tilde{W}_1 \tilde{R}|
= \log |I + P_2 W_1 P_2 R^*| = C(R^*)
$$

(115)

On the other hand, let $R' = U_{2+} \tilde{R} U_{2+}$ and observe that

$$
P_T \geq \text{tr}(R^*) \geq \text{tr}(\tilde{R}')
$$

(116)
so that $R'$ is feasible for (P1) and hence

$$C(R') \geq C(R') = \log |I + \hat{W}_{1}R'\hat{R}| = \tilde{C}(\hat{R})$$

(117)

and finally $C(R') = \tilde{C}(\hat{R}), \mu_1 = 0$ (since the TPC is redundant for the original problem and hence for both problems) and the desired result follows.

Note that Proposition 16 establishes the optimality of projecting all matrices on the sub-space $R(W_2)$ and solving the projected problem instead, if the TPC is not active and $W_2$ is rank-deficient, i.e. if $W_\mu$ is rank-deficient.

Adopting the KKT condition in (81) to the problem in (110), one obtains:

$$(I + \hat{W}_{1}R')^{-1}\hat{W}_{1}R' = \mu_2\tilde{\Lambda}_2\hat{R}$$

(118)

so that

$$r(\hat{R}) = r(\tilde{\Lambda}_2\hat{R}) = r(W_1\hat{R})$$

$$\leq \min(r(W_1), Xr(\hat{R})) \leq r(W_1) \leq r(\hat{R})$$

(119)

and, from (111), $r(R') = r(\hat{R})$, so that $r(R') \leq r(W_1)$, as desired.

E. Proof of Proposition 11

To establish these results, we need the following technical Lemma, which can be established via the standard continuity argument.

**Lemma 1:** Let $W = \lambda uu^T$ be rank-one positive semi-definite matrix, $\lambda > 0$. Then,

$$(I - W^{-1})_{+} = (1 - \lambda^{-1} + uu^T)$$

(120)

Note that the $(\cdot)_{+}$ operator eliminates all singular modes of $W$ and hence its singularity is not a problem, which is somewhat similar to using pseudo-inverse for a singular matrix.

To prove the 1st case, we assume that $W_2$ is not singular and discuss the singular case later. Setting $W = W_2^{-\frac{1}{2}}W_1W_2^{-\frac{1}{2}}$ and applying this Lemma to $(I - W_2^{-\frac{1}{2}}W_1W_2^{-\frac{1}{2}})^{-1}$ in (10), one obtains $R'$ as in (45), after some manipulations, with $W_2 = W_2^{-\frac{1}{2}}$. The condition $\gamma_1 < \gamma_2$ ensures that the TPC is redundant, so that $\mu_1 = 0$ and hence $W_2\mu = \mu_2W_2 > 0$, $tr(W_2R') = P_1$ (since the IPC is active).

If $W_2$ is singular and the TPC is redundant, then one can project all matrices on $R(W_2)$ and solve the projected problem instead without loss of optimality, as was shown in Proposition 16. After some manipulations, this can be shown to result in using the pseudo-inverse instead of the inverse of $W_2$.

To prove the $\gamma_1 \geq \gamma_2$ case, note that, under this condition, $R' = P_1u_1u_1^T$ is feasible under the joint constraint (TPC+IPC). Since it is also optimal without the IPC, it has to be optimal under the joint constraints as well. This proves the “if” part. To prove the “only if” (necessary) part, observe that if $P_1u_1W_2u_1 > P_1$, then $R' = P_1u_1u_1^T$ is not feasible and hence cannot be optimal under the IPC.

To prove the last case, $\gamma_1 \leq \gamma_2 < \gamma_2$, use (10) and note that both constraints are now active (since neither (45) nor $R' = P_1u_1u_1^T$ are feasible under the stated conditions). Applying Lemma 1 as in the 1st case, one obtains (47) after some manipulations.

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