LIST SCHEDULING ALGORITHMS CLASSIFICATION: AN ANALYTICAL STUDY

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Abstract: Multiprocessor scheduling is another name of task scheduling in form of algorithms that are mostly utilized in systematic as well as engineering appliance that is also known as the issue of NP-complete. The main aim of scheduling is the reduction of execution time. The illustration of task scheduling for multiprocessor scheduling is shown by DAG (Directed Acyclic graph). The categorization for this is into Static as well as dynamic scheduling. The list task scheduling is the example of static task scheduling algorithm. Varied task scheduling algorithms, like ISH, HLFET, MCP, ETF, CNPT and DLS are reviewed in this paper. The comparison of list task scheduling is dependent on metrics, termed as SLR, load balancing, efficiency and speed up.

Keywords: Parallel processing, list scheduling, DAG (Directed Acyclic Graph), heuristic based algorithm

I. INTRODUCTION

There is a lot of demand for high speed computing in different application areas. Because of the increasing power of computing, more and more power search has started increasing. A huge issue couldn’t be explained by means of sequential machine in some span of time; therefore, a parallel machine with number of issues has been categorized and has been assign to associate problems on separate processors [1].

Parallel computing is most effective method for fulfilling the calculative constraint for different engineering and scientific applications. Two causes are there for the attractiveness and popularity of parallel computing:
- Lessing computer hardware cost
- Application performance that cannot be executed by conservative computers

Task allocation for the accessible processors for precedence limitation amid tasks for the processors is the aim of task scheduling. In this, the execution time should be less. The representation of application plan is characterized by DAG (Directed Acyclic Graph). Task scheduling is known as NP-completed for some of the restricted cases [2]. Two variants are there, which are connected with task scheduling algorithm. Initial variant is either for judging the presence of communication time and the subsequent variant can be multiprocessor or homogenous system. It is consisted of three components:
- Performance of Homogenous processor
- Task mapping on processors
- Execution sequence for task on every processor.

II. DAG (DIRECTED ACYCLIC GRAPH)

DAG is also termed as TG (Task-Graph) and be considered as general model for parallel program composed of vertices and edges set. DAG is consisted of four tuples, G= {V,E,W,D} in which every vertex, that is V={T1,T2,…, Tn} is taken as graph’s task. E as directed edges seteij depicting the two tasks dependency, Ti plus Tj. Every graph edge is connected by weight D, known as communication time/communication cost [3].

![DAG model having six tasks](image-url)
The representation of comm. of unification time is $D(T_i, T_j)$. The mapping of task set on set $P= (P_1, P_2, \ldots, P_p)$ for $p$ processors and every task can be implemented on processor termed as computation time $W$ and is termed as $W(T_i)$ [4].

Communication time among two tasks is considered as zero on similar processor. Now, the precedence constraint constantly holds $T_j$ that cannot be implemented till $T_i$ that finishes when $i<j$. If the link among $T_i$ and $T_j$ is direct than $T_j$ is considered as the $T_i$ successor which is the $T_j$ predecessor. DAG Illustration has six nodes as illustrated in above figure 1 [5].

An entry task is the task in which there are no predecessors and the exit task is the task in which there are no successors. $T_1$ is known as the entry task and $T_6$ is the exit task.

Communication cost $C(T_i, T_j)$ and computation time for specified DAG is defined below [6]:

| $W(T_1)$ | $W(T_1)$ | $W(T_3)$ |
|---------|---------|---------|
| $W(T_4)$ | $W(T_5)$=2 | $W(T_6)$=3 |
| $C(T_1, T_2)$=5 | $C(T_1, T_3)$=4 | $C(T_2, T_4)$=3 |
| $C(T_2, T_5)$=4 | $C(T_3, T_6)$=5 | $C(T_4, T_6)$=5 |
| $C(T_5, T_6)$=2 |

III. List Task Scheduling Algorithms Classification

The classification of task scheduling is in two categories:

i. Deterministic task scheduling

ii. Non-deterministic task scheduling

Deterministic task scheduling is termed as “Static scheduling/compile time scheduling algorithm”. The tasks information, like, task computation time and communication time with the tasks precedence constraints are considered in this type [8].

Non-deterministic tasks are also termed as Run-time scheduling/dynamic scheduling. For this, the task information is depicted only at the execution time [9].

Below figure is illustrating the classification of list task scheduling algorithms which came under heuristic dependent algorithms.

![Figure 2: List scheduling algorithms classification](image-url)
The task scheduling algorithms are categorized in deterministic and non-deterministic algorithms. The deterministic algorithm is divided into heuristic with guided random search dependent algorithms. The heuristic dependent algorithm provides enhanced solution and satisfactory performance with the polynomial time complexity as compared to ‘exponential time complexity’. It is further divided as: List task scheduling, clustering plus task duplication based algorithms [10].

‘List task scheduling algorithms’ is known as scheduling algorithms in which the assigning of tasks in done by DAG provided to processors. It is easy and provides less complexity as contrast to another algorithm. The clustering algorithm lessens the communication time among DAG tasks. It is an effective algorithm. It executes in two phases:

i. Task clustering

ii. Post clustering

Task duplication based algorithm provides better effectiveness and less scheduling length because of the reduction of communication time with the tasks. It assists for lessening the initialization time for the tasks that are waiting. These algorithms have an aim for processors usage in accurate time [11].

### Table II. Task scheduling algorithms

| Task scheduling algorithms                        | Description                                                                 |
|--------------------------------------------------|-----------------------------------------------------------------------------|
| HLFET (Highest level first with estimate time)   | It is simple and known as the first algorithm of ‘list task scheduling algorithm’. Task priority is chosen by the attributes of static levels. Consideration of communication is not taken place. |
| ISH (Insertion scheduling heuristic) algorithm   | It is effective as it uses the appropriate time being developed by incomplete schedule on processors. The priority of task has been done with static b-level attribute. Provision of enhanced results in contrast of HLFET algorithm is considered in this. |
| MCP (Modified Critical path) algorithm           | It finds task priority by utilizing ALAP (At last as possible) attribute. It gives more priority to tasks that takes less start time. The main con is that it doesn’t have communication time for task priority. |
| ETF (Earliest time first) algorithm              | It finds the earliest start time for each task and later chooses the task having less initial time. Main limitation is that it reduces the scheduling length on each level |
| DLS (Dynamic level scheduling) algorithm         | It finds the task priority on the tasks priority on dynamic basis. It is same as ETF algorithm but DLS utilizes DL attribute while ETF utilizes static level attribute. It doesn’t sustain scheduling list on scheduling procedure. |
| CNPT (Critical node parent tree) algorithm       | It achieves more accuracy and reduces complexity. The prioritization of task is determined with CN (Critical node) attribute. It has two stages; Listing plus Processor assigning phase. It has better performance as contrast to DLS, MCP plus ETF algorithms. |

Below tables shows the time complexity of task scheduling algorithms and priority attributes of task scheduling algorithms [12]. The algorithms considered are HLFET, ISH, MCP, ETF, DLS plus CNPT.

### Table III. Task scheduling algorithms time complexity

| S. No. | Algorithms | Complexity |
|--------|------------|------------|
| 1      | HLFET      | O(v^2)     |
| 2      | ISH        | O(v^2)     |
| 3      | MCP        | O(\(v^2\log v)+p\) |
| 4      | ETF        | O(pv^3)    |
| 5      | DLS        | O(pv^3)    |
| 6      | CNPT       | O(v^2)     |

### Table IV. Priority Attribute of task scheduling algorithms

| S. No. | Algorithms | Priority attribute |
|--------|------------|--------------------|
| 1      | HLFET      | Static-level       |
| 2      | ISH        | Static-level       |
| 3      | MCP        | ALAP               |
| 4      | ETF        | Static-level       |
| 5      | DLS        | Dynamic-level      |
| 6      | CNPT       | Critical-node      |

### IV. COMPARATIVE METRICS

The analysis of performance could be executed in list task scheduling algorithms on the basis of comparison metrics, SLR (Scheduling length ratio), Speed up, Efficiency and Load balancing [13-14].
Table V. Task scheduling comparative metrics

| Metrics                  | Description                                                                 |
|--------------------------|------------------------------------------------------------------------------|
| SLR (Scheduling length ratio) | It is the proportion of scheduling length to average execution time on each processor. Mathematically, it can be described as: $\text{efficiency} = \frac{S_p}{N_p}$ As shown, $S_p$ is the speed up and $N_p$ is the number of processors. |
| Speed up                 | It is the proportion among ‘sequential execution time’ with ‘parallel execution time’ in which the sequential execution time as the amount of total computation time for every task with ‘parallel execution time’ which is the SL on less amount of processors. $\text{Speedup} = \frac{\sum_{i=1}^{n} T_i}{T_p}$ As shown, $\sum_{i=1}^{n} T_i$ is the amount of computational task in chronological order as $1,2,3,\ldots,n$ and $T_p$ is schedule length and total parallel execution time of DAG |
| Efficiency               | It is the measurement of processor utilization. Mathematical, it can be described as: $\text{efficiency} = \frac{S_p}{N_p}$ As shown, $S_p$ the speed up and $N_p$ is the number of processors |
| Load balancing           | It can be described with the proportion of scheduling length to average execution time on each processor. |

V. RELATED WORK

This section describes the work till date for different scheduling techniques for fulfilling varied QoS (Quality of service) metrics with energy saving methods.

Table VI. A glance of existing techniques

| Author                  | Proposed work                                                                 | Research Gap                                                                 |
|-------------------------|------------------------------------------------------------------------------|------------------------------------------------------------------------------|
| Abdul Razzaque et al. [1]| Task Scheduling in Cloud Computing on the basis of carbon footprint         | It has smaller amount flexibility and less reliability approach for more execution time. |
| Hao Wu et al. [2]       | Optimization of Deadline Constrained DAG Applications by using VM (Virtual machine) concept. | The accessible task model is relied on task's execution times be hard for efficiently calculating in a cloud environment. Other issue is the cloud environment virtualization overhead for proposed algorithm that could be improved on the virtualization overhead basis. |
| Amandeep Verma et al. [3]| Cost and Time aware Scheduling strategy for Workflow application being executed in Cloud | Because of less optimal schedule plan for real cloud environment, the computational cost is additional and there is a possibility of enhancement in the enhanced schedule planning. |
| Hamid Arabnejad et al. [4]| Scheduling Algorithm on the basis of budget constraint provided by user for Workflow Applications | Because of requirement of dynamic concurrent DAG scheduling issue, that execute parallel workflows which cannot execute together but might distribute the resources for total cost for the user that could be lessen |
| Jia Yu and R. Buyya [5]  | DAG dependent scheduling for budget constraint satisfaction by metaheuristic genetic algorithm on efficacy grids | Usage of optimization method has not considered for solving QoS constraints for security and reliability. |
| Wei Zheng et al. [6]    | 'Budget-Deadline Constrained Workflow Planning used for Admission Control for Bi-criteria DAG scheduling' | The computational complexity of presented work is higher so the success rate is reduced and in presented work, the middle DAG scheduling heuristic technique is required. |
| Zhuo Tang et al. [7]    | DVFS enable energy effective workflow task scheduling                       | Power consumption lessens by 46.5% but slacked makespanenhances. |
| Weihong Chen et al. [8] | Effective Task Scheduling used for Budget Constrained Parallel appliance on Heterogeneous Cloud Computing scheme | The issue of proposed work is only appropriate for the homogeneous cloud environment. |
| Jasraji Meena et al. [9] | Cost Effective GA in favor of Workflow Scheduling in Cloud in Deadline Constraint | There is a big issue of shutdown time of VMs and due to the general execution workflow cost is affected. Due to the absence of optimal schedule plan for a real cloud environment, the computational cost is more and there is a chance of improvisation in the optimal schedule planning. |
| Anton Beloglazov et al. [10]| Heuristics of energy awareness in resource allocation for effective data center management | Difficult to run on large-scale and at large-scale energy consumption is more. There is no any concept of the generic resource manager. |

VI. CONCLUSION

The list task scheduling algorithms classification HLFET, ISH, MCP, DLS, ETF, and CNPT algorithms which are of homogenous environment has been studied and analyzed in this paper. The pros and cons of time complexity are considered. The algorithms are dependent on few priority attributes. As per priority attributes, the assigning of priority is taken place. Some computation metrics, namely, SLR, load balancing, efficiency and speed up have been studied that
provides the assistance for differentiating the algorithms. The list scheduling algorithms provides more effectiveness and less SL than another scheduling algorithm.

VII. REFERENCES

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