Blockchain Based Packet Delivery Mechanism for WSN

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Abstract: The latest trend in the research filed implies the importance of data security in wireless sensor networks. There are various approaches identified for securing the data by using trust wide security such as cryptographic systems and routing protocols. However, these approaches are very critical to identify the optimal path in the network and attacks by an unauthorized node cannot be prevented. In this paper, a new algorithm for combining the AODV (Ad Hoc On-Demand Distance Vector) routing protocol and particle swarm optimization (PSO) is implemented to produce a trustable routing in every location through block chain. The possible routing procedure will enhance the routing nodes to acquire routing information among all the nodes but will never allow the node to capture the information. The routing protocol on the blockchain is used to utilize the path efficiently without deviation caused by other anchor nodes. It also identifies the congestion in the entire path of the particular network and avoids tampering of information between the nodes. The blockchain enabled with PSO algorithm and AODV routing protocol provides the simulation results about the efficient packet delivery system. The Security has been performed in every node used to identify the best route for producing the efficient throughput and quality of services.

Index Terms: Wireless sensor networks (WSN), PSO, AODV, Blockchain

I. INTRODUCTION

A wireless sensor network (WSN) with autonomous sensors monitors the physical or environmental conditions, such as temperature, sound, pressure, etc. and also cooperatively pass their data through the network to the prime location [1]. The extensive deployment of sensor networks is quite practical these days. A network of this numerous sensors allows an efficient solution to various challenging tasks: traffic monitoring, monitoring of building with respect to the structure, fire, and security, military sensing and tracking, distributed measurement of seismic activity, real-time pollution monitoring, wildlife monitoring and wildfire tracking [2]. The existing key pre-distribution schemes such as basic probabilistic and q composite schemes used in the sensor networks forms a pair wise key for the establishment and authentication between sensor nodes and mobile sinks. An attacker places a replicated mobile sink preloaded with some compromised keys to obtain a large number of keys by capturing a small fraction of nodes to gain control of the network [3]. To overcome this, a three-tier security scheme was further proposed that makes use of the polynomial pool based scheme [4] and pairwise authentication schemes, which further reduces these attacks to some extent but the energy level in this scheme is not high. Mobile ad hoc network (MANETs) is a wireless network that has a routable networking environment on top of a link layer in ad hoc network. MANET is a peer-to-peer self-forming and self-healing network which has a central controller to determine, optimize, and distribute the routing table which is not present in mesh network. The device of the MANET has the power to move independently in any direction with many anonymous MANET routing protocols like AODV or DSR [6-9]. The different cryptosystems are used to implement the same function to measure the performance of routing changes. It is an added advantage for ANODR to overcome the inefficient anonymous routing of MANET. An improved AODV routing protocol which is based on minimal route, is used to simulate the throughput and reduce the end-to-end delay and other parameters for evaluating the performance of the wireless sensor network by OPNET[11]. The destination sequence distance vector (DSDV) and ad hoc on demand distance vector (AODV) routing protocols are popularly used in mobile ad hoc networks. Among the two protocols, AODV is more reliable with the affordable burden compared to DSDV [12]. The performance evaluation of different network parameters on different topologies based on varying the pause time with respect to constant speed (node speed) in different terrain areas as small (1000 m. x 1000 m.), large (2000 m. x 1000 m.) and very large (2000 m. x 2000 m.) using AODV routing protocol and monitoring of critical conditions with the help of important parameters like packet delivery fraction, average end-to-end delay, average throughput, NTL and packet loss [13]. Energy Aware Routing Protocol (ETARP) is used to discover and select routes based on the maximum utility to incur additional cost in overhead compare to the common AODV routing protocol [14]. The flip ambiguity problem has been solved by distributed two-phase PSO algorithm to improve the efficiency and precision of the entire network. Moreover, the unknown nodes with minimum two or three near-collinear references are made to be localized in our research [15]. A multi-objective particle swarm optimization algorithm was used to find the optimal solution. The global optimum is obtained according to the proportion of selection [16]. The proof of work (PoW) in original bitcoins is to carry the pulling out and to store new data blocks by blockchain technology. It also adopts the reliable data possession to replace,[18]. A novel blockchain-based contractual routing (BCR) protocol is used for a network of untrusted IoT devices. In view of the traditional secure routing protocols in which a central authority (CA) is required to formulate the identification and authentication of each device. It produced the routing overhead of the BCR protocol in 5 times lower compared to AODV at the cost of a slightly lower packet delivery ratio. BCR is fairly resistant to both Blackhole.
and Greyhole attacks. The results show that the BCR protocol enables distributed routing in different IoT networks.[19]. Reinforcing the trust in distributed environments without the need of authorities is a technological advance which will change the industry’s perspective. The IoT which is one among the disruptive technologies such as big data and cloud computing. To overcome its limitations since its conception, blockchain will be one of the next replacements [20]. The feasible routing information is given by routing scheme obtaining the routing nodes on the blockchain which will make the routing information traceable and impossible to tamper with [21].

II. ARCHITECTURAL MODEL OF AODV AND PSO

Figure 1 shows the architecture model of congestion avoiding flow system of the routing protocol. The AODV protocol has been used to send the packet from source to destination. If either source or destination has been attacked by a malicious node for affecting all the packets, the PSO algorithm used to gather the authorized node by using the method of blockchain. After the arrival of all the packets in every node, overall performance has been analyzed with statistical measurement of poison distributed scheme. This scheme used to measure the performance of the authorized node for obtaining the smooth flow control of the packets in every route. Figure 1 shows N1, N2, Nn are a number of nodes. Each node is sending the pre-request (PREQ) as a packet to the destination. The receiver has observed each path and finding the packets without a malicious node. The receiver then sends the reply as an acknowledgment to the source. In this case, the packet delay time has been identified and sending the packet without any caution of delay.

III. METHODOLOGY OF WORKING PRINCIPLES OF AODV AND PSO

The AODV has used to implement in various networks to adopt the route efficiently with same energy level by using PSO on blockchain technology. The reason behind avoiding congestion is used:

- To increase the efficient throughput of the system.
- To increase the lifetime of the network.
- To maintain the quality of services.

Figure 2 indicates the sequence numbers for making node communication within the time stamps. The nodes can compare themselves about the priority of their information among nodes. The sequence number increases with respect to the type of message send by each node. This sequence number of each node is shared to rest of the nodes.

Particle Swarm Optimization (PSO) is a heuristic designed algorithm which creates no assumptions about the problem being optimized. PSO can be used in candidate solutions on optimization problems that are partially irregular, noisy, change over time, etc[17].

IV. WHY BLOCKCHAIN REQUIRED ON AODV AND PSO?

The blockchain is an imperishable digital ledger of economic transactions that can record all financial transactions with its virtual values. The blockchain network has decentralized authority to communicate the information between the nodes. Also, it is a shared and absolute ledger, where the information is open for all nodes. When any intrusion takes place by a node within the network, the other nodes get the prompt information about this malfunction.
Therefore, the overall actions that is built on the blockchain is transparent and everyone involved is accountable for their actions [22]. In this connection, the blockchain technology has been implemented in every node. If the nodes are used to capture the route to its destination, every node blocks the shortest route to reach its recipient point. These situations arise that every node can be visibly digesting every other node information, but cannot be theft and modify the information about another node. The PSO algorithm is used only for measuring the route which is not affected by malicious users. In this case of blockchain technology, the route can be blocked by every node which in turn only authorized node can be blocked as a chain and information might be shared only authorized node. The malicious node cannot be interfered in any situation to control every data. In order to obtain the blockchain technology in AODV, produced the efficient energy level and improve the throughput for producing the Quality of Service.

V. PACKET DELIVERY RATIO CALCULATED USING POISON DISTRIBUTION

A discrete random variable X with assumed values X = 0, 1, 2, 3,..∞ is said to follow Poisson distribution if its probability mass function is

\[ P(X=x) = \begin{cases} \frac{e^{-\lambda} \lambda^x}{x} , & x = 0, 1, 2, ..\infty; \text{o, otherwise} \\ \end{cases} \]  

(1)

Here \( \lambda \) is the parameter of the Poisson distribution. Let X denotes the number of nodes \( x_1, x_2...x_i \) in every direction of network during packet transmission. The nodes establish the communication through their anchor nodes in the form of \( x_1->x_2, x_2->x_3, x_3->x_i \). The AODV protocol is used to measure every path for obtaining the request from the source and it is used to connect to destination through every hop. The packets are transferred irrespective of the number of hops already connected. When node fails due to the network disruption it gets retransmitted from the source. These cases analyze through PSO technique to optimize the congestion path where it is exactly occurred. Sometimes the malicious nodes try to interrupt communication through accumulating the paths between the nodes. Hence the nodes are decentralized to communicate each other and control by blockchain. The properties like time and information of the nodes are shared among the authorized nodes alone. The probability of infinite number of events is measured using the Poisson distribution. This calculation gives the number of nodes successfully linked to reach its destination path without any congestion in stipulated time. Equation 2 shows the estimated process time of packet delivery irrespective of a number of nodes.

\[ \text{DelayTime} = \frac{\text{ProcessTime}}{\text{EstimationTime}} X \cdot \text{No. of Nodes} \]  

(2)

VI. ALGORITHM:

INTEGRATION OF AODV AND PSO

Step1: Initialize number of nodes \( n_1, n_2, n_3, n_4, ..., n_i \)
Step2: Calculate a number of Hops Let \( h=\text{hopCount} \)
Step3: Initialize the router=\( i \) where \( i=1, 2, 3, ...k \)
Step4: Identifying the position of every path with respect to a number of nodes \( n_1, n_2, ..., n_i \)
Step5: if \( (i>1) \)
    \( \text{if}((n_1+i)+(n_2+i) < \text{Previous estimate router}) \)
    i) Calculate Transmission rate(\( T_r \))
    \( \text{where} \ T_r=\text{Time of Node} \ n_1 \rightarrow \text{Time of Node} \ n_2 \)
    Let \( \text{N}_1=n_1+i \rightarrow \text{Starting node of the route} \)
    Let \( \text{N}_2=n_2+i \rightarrow \text{Neighbor node of the route} \)
    Similarly calculate \( \text{N}_3, \text{N}_4, ..., \text{N}_i \)/Link the nodes \( \text{N}_1, \text{N}_2, ..., \text{N}_i \)

Calculate Distance= \( \sqrt{(N_i-N_j)^2+(N_i-N_j)^2} \)

//Find the Distance between nodes to identify //shortest distance
Calculate Process Time= \( \frac{\text{Distance}}{T_r} \)

\[ \text{DelayTime} = \frac{\text{ProcessTime}}{\text{EstimationTime}} \cdot X \cdot \text{No. of Nodes} \]

//Where (i) Process Time-> Starting time of the node
// (ii) Estimation time-> Ending time of the node
Link to router i with neighbor node to set short distance based on Delay Time.
Increase the hop count \( h_{++} \) and node is connected to neighbor node i.e. \( n_{++} \)
else

Remove the Unused link from the router.
the router is not initialized.

Step6: To find the malicious node by using blockchain
For \( h \) in 0 to \( n \)
Let block=0; target=n.
\text{if} \ (\text{distance}<\text{target})
    The current block is moved to the next block for the authorized link.
else

Step7: End the process of an Algorithm
VII. RESULT AND SIMULATION

The simulation result of this research work has been shown how the AODV and PSO algorithm can be executed using the probabilistic approach of Poison distribution. Equation 1 shows the poison distribution value of \( n=10 \) and the estimated time has been measured in terms of AODV. The average time delay is measured in the point of view is 0.17. Each node has to travel around the other node with time delay. In this scenario 0.17 is the ratio has been done by AODV protocol. Figure 3 shows the average time delay has happened from node 1 to 10.

![Figure 3: Probabilistic value of time delay using poison distribution](image)

Table 1 shows the delay time calculation for packet delivery with process time in seconds and estimated time. This calculation is based on equation 2 for observing the utilization of the communication path. Figure 4 shows the delay time of packet delivery between source and destination based on the calculation of estimated time in packet delivery. The average delay time of packet transmission with \( n=5 \) (no of nodes) and Estimated time of 5 nodes is 0.999. In every hop, the processing time and the estimated time has been measured. The processing time has awaited the communication of node 0.03 in seconds for 5 nodes. That is every hop of the distance is varying the certain level and finally reaches to the destination with 0.03 in seconds.

![Figure 4: Measurement of Delay time between Estimated and process time](image)

Table 1: Calculation of Process Time and Estimated Time

| No of Nodes | Process Time (in seconds) | Estimated Time (in seconds) |
|-------------|---------------------------|-----------------------------|
| 0           | 0.368                     | 0.368                       |
| 1           | 0.368                     | 0.736                       |
| 2           | 0.184                     | 0.919                       |
| 3           | 0.061                     | 0.981                       |
| 4           | 0.015                     | 0.996                       |
| 5           | 0.03                      | 0.999                       |

When the packet is to send from source to destination, the distance and time have been measured. If average time is exceeded the estimated time, then the PSO algorithm used to identify the congestion and where exactly the path has to be connected. Table 2 shows the Pre-request (PREQ) of packet delivery from source to destination. The number of packets has delivered with data 0, 1…10 and the delay time is measured in every path. The final measurement was taken as 99.16 for 50 seconds is the approximate time for more data with quality of service. Figure 5 shows how the data had been flowing from source to destination that represents in the graphical picture. The flow of data and time delay in every path are represented in figure 5. Once the packet delivery ration has identified by using AODV and PSO, then the security measures should be implemented by using blockchain. The data has been authorized in terms of security. The information is distributed in the intermediate nodes, but an only authorized node can communicate and established the communication.

Table 2: Experimental Results of RREQ for measuring packet delivery in AODV

| No of Nodes | Pause Time (m/s) | Packet Delivery Ratio |
|-------------|------------------|-----------------------|
| 1           | 0                | 99.5                  |
| 2           | 10               | 99                    |
| 3           | 20               | 99.4                  |
| 4           | 30               | 99.5                  |
| 5           | 40               | 98.76                 |
| 6           | 50               | 99.16                 |

Figure 6 shows how the packet delivery ratios are measured in every node and how it can be transferred from source to destination by
In this paper, we have identified how AODV has integrated with the PSO algorithm for finding the exact delay. The congestion has occurred anywhere of the path, but both algorithms have used to find the exact path for avoiding the congestion. The probabilistic approach of poison distribution has implemented for measuring the delay time when the number nodes are communicated with each other. The efficiency of the AODV protocol has to be identified when the nodes are interconnected by blockchain with the PSO algorithm. The Quality of Service (QoS) has been identified in this paper based on packet delivered ratio from source to destination without any congestion. In this scenario packet delivery ratio has implemented with 99.16 for 100 packets which is accumulated in 50 m/s. Finally, the variation of the packet delivery has been achieved within the estimated time. The future enhancement of the work will be extracted for how the data will be secured using consensus mechanism of blockchain in the various levels of networks using n number of nodes.

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