FACE VECTORS OF FLAG COMPLEXES

ANDY FROHMADER

Abstract. A conjecture of Kalai and Eckhoff that the face vector of an arbitrary flag complex is also the face vector of some particular balanced complex is verified.

1. Introduction

We begin by introducing the main result. Precise definitions and statements of some related theorems are deferred to later sections.

The main object of our study is the class of flag complexes. A simplicial complex is a flag complex if all of its minimal non-faces are two element sets. Equivalently, if all of the edges of a potential face of a flag complex are in the complex, then that face must also be in the complex.

Flag complexes are closely related to graphs. Given a graph $G$, define its clique complex $C = C(G)$ as the simplicial complex whose vertex set is the vertex set of $G$, and whose faces are the cliques of $G$. The clique complex of any graph is itself a flag complex, as for a subset of vertices of a graph to not form a clique, two of them must not form an edge. Conversely, any flag complex is the clique complex of its 1-skeleton.

The Kruskal-Katona theorem \cite{6,5} classifies the face vectors of simplicial complexes as being precisely the integer vectors whose coordinates satisfy some particular bounds. The graphs of the “rev-lex” complexes which attain these bounds invariably have a clique on all but one of the vertices of the complex, and sometimes even on all of the vertices.

Since the bounds of the Kruskal-Katona theorem hold for all simplicial complexes, they must in particular hold for flag complexes. We might expect that flag complexes which do not have a face on most of the vertices of the complex will not come that close to attaining the bounds of the Kruskal-Katona theorem.

One way to force tighter bounds on face numbers is by requiring the graph of the complex to have a chromatic number much smaller than the number of vertices. The face vectors of simplicial complexes of a given chromatic number were classified by Frankl, Füredi, and Kalai \cite{4}.

Kalai (unpublished; see \cite{8} p. 100) and Eckhoff \cite{1} independently conjectured that if the largest face of a flag complex contains $r$ vertices, then it must satisfy the known bounds (see \cite{4}) for complexes of chromatic number $r$, even though the flag complex may have chromatic number much larger than $r$. We prove their conjecture.

Theorem 1.1. For any flag complex $C$, there is a balanced complex $C'$ with the same face vector as $C$. 

Date: May 8.
Our proof is constructive. The Frankl-Füredi-Kalai theorem states that an integer vector is the face vector of a balanced complex if and only if it is the face vector of a colored “rev-lex” complex. This happens if and only if it satisfies certain bounds on consecutive face numbers. Given a flag complex, for each $i$, we construct a colored “rev-lex” complex with the same number of $i$-faces and $(i+1)$-faces as the flag complex, thus showing that all the bounds are satisfied.

The structure of the paper is as follows. Section 2 contains basic facts and definitions related to simplicial complexes. In Section 3, we discuss the Kruskal-Katona theorem and the Frankl-Füredi-Kalai theorem, and lay the foundation for our proof. Finally, Section 4 gives our proof of the Kalai-Eckhoff conjecture.

2. Preliminaries on simplicial complexes

In this section, we discuss some basic definitions related to simplicial complexes.

Recall that a simplicial complex $\Delta$ on a vertex set $V$ is a collection of subsets of $V$ such that, (i) for every $v \in V$, $\{v\} \in \Delta$ and (ii) for every $B \in \Delta$, if $A \subset B$, then $A \in \Delta$. The elements of $\Delta$ are called faces. The maximal faces (under inclusion) are called facets.

For a face $F$ of a simplicial complex $\Delta$, the dimension of $F$ is defined as $\dim F = |F| - 1$. The dimension of $\Delta$, $\dim \Delta$, is defined as the maximum dimension of the faces of $\Delta$. A complex $\Delta$ is pure if all of its facets are of the same dimension.

The $i$-skeleton of a simplicial complex $\Delta$ is the collection of all faces of $\Delta$ of dimension $\leq i$. In particular, the 1-skeleton of $\Delta$ is its underlying graph.

It is sometimes useful in inductive proofs to consider certain subcomplexes of a given simplicial complex, such as its links.

Definition 2.1. Let $\Delta$ be a simplicial complex and $F \in \Delta$. The link of $F$, $\text{lk}_\Delta(F)$, is defined as $\text{lk}_\Delta(F) := \{ G \in \Delta \mid F \cap G = \emptyset, F \cup G \in \Delta \}$.

The link of a face of a simplicial complex is itself a simplicial complex. It will be convenient to define the notion of a link of a vertex of a graph.

Definition 2.2. The link of a vertex $v$ in a graph $G$, denoted $\text{lk}_G(v)$, is the induced subgraph of $G$ on all vertices adjacent to $v$.

Note that $\text{lk}_G(v)$ coincides with the 1-skeleton of the link of $v$ in the clique complex of $G$.

Next we discuss a special class of simplicial complexes known as flag complexes.

Definition 2.3. A simplicial complex $\Delta$ on a vertex set $V$ is a flag complex if all of its minimal non-faces are two element sets. A non-face of $\Delta$ is a subset $A \subseteq V$ such that $A \notin \Delta$. A non-face $A$ is minimal if, for all proper subsets $B \subset A$, $B \in \Delta$.

In the following, we refer to the chromatic number of a simplicial complex as the chromatic number of its 1-skeleton in the usual graph theoretic sense.

We also need the notion of a balanced complex, as introduced and studied in [4].

Definition 2.4. A simplicial complex $\Delta$ of dimension $d-1$ is balanced if it has chromatic number $d$.

Note that the chromatic number of a simplicial complex of dimension $d-1$ must be at least $d$, as it has some face with $d$ vertices, all of which are adjacent, so
coloring that face takes \( d \) colors. A balanced complex is then one whose chromatic number is no larger than it has to be.

Not all simplicial complexes are balanced complexes. For example, a pentagon (five vertices, five edges, and one empty face) is not a balanced complex, because it has chromatic number three but dimension only one.

In this paper, we study the face numbers of flag complexes.

**Definition 2.5.** The \( i \)-th face number of a simplicial complex \( C \), denoted \( c_i(C) \) is the number of faces in \( C \) containing \( i \) vertices. These are also called \( i \)-faces of \( C \). If \( \dim C = d - 1 \), the face vector of \( C \) is the vector

\[
c(C) = (c_0(C), c_1(C), \ldots, c_d(C)).
\]

In particular, for any non-empty complex \( C \), we have \( c_0(C) = 1 \), as there is a unique empty set of vertices, and it is a face of \( C \).

Since flag complexes are the same as clique complexes of graphs, it is sometimes convenient to talk about face numbers in the language of graphs.

**Definition 2.6.** The \( i \)-th face number of a graph is the \( i \)-th face number of its clique complex. Likewise, the clique vector of a graph is the face vector of its clique complex.

The face numbers defined here are shifted by one from what is often used for simplicial complexes. This is done because we are primarily concerned with flag complexes, or equivalently, clique complexes of graphs, where it is more natural to index \( i \) as the number of vertices in a clique of the graph, following Eckhoff [3].

The graph concept corresponding to the dimension of a simplicial complex is the clique number.

**Definition 2.7.** The clique number of a graph is the number of vertices in its largest clique.

Note that the clique number of a graph is one larger than the dimension of its clique complex.

3. The Kruskal-Katona and Frankl-Füredi-Kalai theorems

For the general case of simplicial complexes, the question of which face vectors are possible is answered by the Kruskal-Katona theorem [6, 5]. Stating the theorem requires the following lemma.

**Lemma 3.1.** Given any positive integers \( m \) and \( k \), there is a unique \( s \) and unique \( n_k > n_{k-1} > \cdots > n_{k-s} \geq k - s > 0 \) such that

\[
m = \binom{n_k}{k} + \binom{n_{k-1}}{k-1} + \cdots + \binom{n_{k-s}}{k-s}.
\]

The representation described in the lemma is called the \( k \)-canonical representation of \( m \).

**Theorem 3.2** (Kruskal-Katona). For a simplicial complex \( C \), let

\[
m = c_k(C) = \binom{n_k}{k} + \binom{n_{k-1}}{k-1} + \cdots + \binom{n_{k-s}}{k-s}.
\]
be the $k$-canonical representation of $m$. Then
\[ c_{k+1}(C) \leq \binom{n_k}{k+1} + \binom{n_{k-1}}{k} + \cdots + \binom{n_{k-s}}{k-s+1}. \]

Furthermore, given a vector $(1, c_1, c_2, \ldots, c_t)$ which satisfies this bound for all $1 \leq k < t$, there is some complex that has this vector as its face vector.

To construct the complexes which demonstrate that the bound of the Kruskal-Katona theorem is attained, we need the reverse-lexicographic (“rev-lex”) order. To define the rev-lex order of $i$-faces of a simplicial complex on $n$ vertices, we start by labelling the vertices $1, 2, \ldots$. Let $\mathbb{N}$ be the natural numbers, let $A$ and $B$ be distinct subsets of $\mathbb{N}$ with $|A| = |B| = i$, and let $A \cap B$ be the symmetric difference of $A$ and $B$.

**Definition 3.3.** For $A, B \subset \mathbb{N}$ with $|A| = |B|$, we say that $A$ precedes $B$ in the rev-lex order if $\max(A \cap B) \in B$, and $B$ precedes $A$ otherwise.

For example, $\{2, 3, 5\}$ precedes $\{1, 4, 5\}$, as $3$ is less than $4$, but $\{3, 4, 5\}$ precedes $\{1, 2, 6\}$.

**Definition 3.4.** The rev-lex complex on $m$ $i$-faces is the pure complex whose facets are the first $m$ $i$-sets possible in rev-lex order. This complex is denoted $C_i(m)$.

We can also specify more than one number in the face vector. For two sequences $i_1 < \cdots < i_r$ and $(m_1, \ldots, m_r)$, let
\[ C = C_{i_1}(m_1) \cup C_{i_2}(m_2) \cup \cdots \cup C_{i_r}(m_r). \]

The standard way to prove the Kruskal-Katona theorem involves showing that if the numbers $m_1, \ldots, m_r$ satisfy the bounds of the theorem, then the complex $C$ has exactly $m_j$ $i_j$-faces for all $j \leq r$ and no more. In this case, we refer to $C$ as the rev-lex complex on $m_1$ $i_1$-faces, $\ldots$, $m_r$ $i_r$-faces.

For example, if the complex $C$ has $\binom{9}{3} + \binom{6}{3} = 99$ 3-faces, then the Kruskal-Katona theorem says that it can have at most $\binom{9}{3} + \binom{6}{4} = 146$ 4-faces. The rev-lex complex on 99 3-faces and 146 4-faces gives an example showing that this bound is attained.

The 1-skeleton of the rev-lex complex that gives the example for the existence part of the Kruskal-Katona theorem always has as large of a clique as is possible without exceeding the number of edges allowed, as well as a chromatic number of either the number of non-isolated vertices or one less than this. It turns out that if we require a much smaller chromatic number, we can get a much smaller bound. To take an extreme example, if $c_3(C) = 1140$, then the Kruskal-Katona theorem requires that $c_4(C) \leq 4845$. But if we require the complex $C$ to be 3-colorable, then we trivially cannot have any faces on four vertices, and $c_4(C) = 0$.

We could ask what face vectors occur for $r$-colorable complexes for a given $r$. This was solved by Frankl, Füredi, and Kalai [4]. In order to explain their result, we need the concept of a Turán graph.

**Definition 3.5.** The Turán graph $T_{n,r}$ is the graph obtained by partitioning $n$ vertices into $r$ parts as evenly as possible, and making two vertices adjacent exactly if they are not in the same part. Define $\binom{n}{k}_r$ to be the number of $k$-cliques of the graph $T_{n,r}$. 
The structure of the Frankl-Füredi-Kalai theorem is similar to that of the Kruskal-Katona theorem, beginning with a canonical representation of the number of faces.

**Lemma 3.6.** Given positive integers \( m, k, \) and \( r \) with \( r \geq k \), there are unique \( s, n_k, n_{k-1}, \ldots, n_{k-s} \) such that

\[
m = \binom{n_k}{k} + \binom{n_{k-1}}{k-1} + \cdots + \binom{n_{k-s}}{k-s},
\]

\[
n_{k-i} - \left\lfloor \frac{n_{k-i-1}}{r-1} \right\rfloor > n_{k-i-1} \quad \text{for all} \quad 0 \leq i < s, \quad \text{and} \quad n_{k-s} \geq k-s > 0.
\]

This expression is called the \((k,r)\)-canonical representation of \( m \).

**Theorem 3.7** (Frankl-Füredi-Kalai). For an \( r \)--colorable complex \( C \), let

\[
m = c_k(C) = \binom{n_k}{k} + \binom{n_{k-1}}{k-1} + \cdots + \binom{n_{k-s}}{k-s},
\]

be the \((k,r)\)-canonical representation of \( m \). Then

\[
c_{k+1}(C) \leq \binom{n_k}{k+1} + \binom{n_{k-1}}{k} + \cdots + \binom{n_{k-s}}{k-s + 1}. 
\]

Furthermore, given a vector \((1, c_1, c_2, \ldots, c_t)\) which satisfies this bound for all \( 1 \leq k < t \), there is some \( r \)--colorable complex that has this vector as its face vector.

The examples which show that this bound is sharp come from a colored equivalent of the rev-lex complexes of the Kruskal-Katona theorem.

**Definition 3.8.** A subset \( A \subset \mathbb{N} \) is \( r \)--permissible if, for any two \( a, b \in A \), \( r \) does not divide \( a - b \). The \( r \)--colored rev-lex complex on \( m \) \( i \)--faces is the pure complex whose facets are the first \( m \) \( r \)--permissible \( i \)--sets in rev-lex order. This complex is denoted \( C^r_i(m) \).

The complex \( C^r_i(n) \) is \( r \)--colorable because we can color all vertices which are \( i \) modulo \( r \) with color \( i \).

As with the uncolored case, we can define a rev-lex complex with specified face numbers of more than one dimension. For two sequences \( i_1 < \cdots < i_s \) and \((m_1, \ldots, m_s)\), let \( C = C^r_{i_1}(m_1) \cup C^r_{i_2}(m_2) \cup \cdots \cup C^r_{i_s}(m_s) \). The proof of Theorem 3.7 involves showing that if the numbers \( m_1, \ldots, m_r \) satisfy the bounds of the theorem, then the complex \( C \) has exactly \( m_j \ i_j \)--faces and no more. In this case, we refer to \( C \) as the \( r \)--colored rev-lex complex on \( m_1 \) \( i_1 \)--faces, \( \ldots, m_r \) \( i_r \)--faces. This complex is likewise \( r \)--colorable with one color for each value modulo \( r \).

In the case of flag complexes, the face numbers of the complex must still follow the bounds imposed by the chromatic number by Theorem 3.7. Still, there are graphs whose clique number is far smaller than the chromatic number, and having no large cliques seems to force tighter restrictions on the clique vector than the chromatic number alone. In particular, given a graph \( G \) of clique number \( n \), we must have \( c_i(G) = 0 \) for all \( i > n \), while the bound from the chromatic number and Theorem 3.7 may be rather large. Note that the chromatic number must be at least the size of the largest clique, as any two vertices in a maximum size clique must have different colors.

It has been conjectured by Kalai (unpublished) and Eckhoff that, given a graph \( G \) with clique number \( r \), there is an \( r \)--colorable complex with exactly the
same face numbers as the clique complex of the graph. Their conjecture generalizes
the classical Turán theorem from graph theory, which states that among all triangle-
free graphs on \( n \) vertices, the Turán graph \( T_{n,2} \) has the most edges \([9]\). The goal of
the following section is to verify Theorem 1.1, proving their conjecture.

4. Proof of the Kalai-Eckhoff conjecture

Fix a graph \( G \) with \( c_{r+1}(G) = 0 \) and fix \( k \geq 0 \). We start by showing that
there is an \( r \)-colorable complex \( C \) with \( c_k(G) = c_k(C) \) and \( c_{k+1}(G) = c_{k+1}(C) \) (see
Lemma 4.1 below).

The case \( k = 1 \) of the lemma is given by Turán’s theorem \([10]\). It was generalized
by Zykov \([10]\) to state that if \( G \) is a graph on \( n \) vertices of chromatic number \( r \),
then \( c_i(G) \leq \binom{n}{i} \). The case \( k = 2 \) was proven by Eckhoff \([2]\). A subsequent paper
of Eckhoff \([5]\) established a bound on \( c_i(G) \) in terms of \( c_2(G) \) for all \( 2 \leq i \). All of
these results are special cases of our Theorem 1.2 and proven independently below.

**Lemma 4.1.** If \( G \) is a graph with \( c_{r+1}(G) = 0 \) and \( k \) is a nonnegative integer, then
there is some \( r \)-colorable complex \( C \) with \( c_k(C) = c_k(G) \) and \( c_{k+1}(C) = c_{k+1}(G) \).

Proof. We use induction on \( k \). For the base case, if \( k = 0 \), take \( C \) to be a
complex with the same number of vertices as \( G \), no edges, and all vertices the same
color.

Otherwise, assume that the lemma holds for \( k - 1 \), and we need to prove it
for \( k \). The approach for this is to use induction on \( c_{k+1}(G) \). For the base case, if
\( c_{k+1}(G) = 0 \), then there is trivially some \( r \)-colorable complex \( C \) with \( c_k(C) = c_k(G) \)
and \( c_{k+1}(C) = 0 \).

For the inductive step, suppose that \( c_{k+1}(G) > 0 \). Let \( v_0 \) be the vertex of \( G 
\) contained in the most cliques of \( k + 1 \) vertices; in case of a tie, arbitrarily pick some
vertex tied for the most to label \( v_0 \). Let the vertices of \( G \) not adjacent to \( v_0 \) be
\( v_1, v_2, \ldots, v_s \).

Given a graph \( G \) and a vertex \( v \), there is a bijection between \( k \)-cliques of \( \text{lkg}(v) \) and
\((k+1)\)-cliques of \( G \) containing \( v \), where a \( k \)-clique of \( \text{lkg}(v) \) corresponds to the
\((k+1)\)-clique of \( G \) containing the \( k \) vertices of the \( k \)-clique of \( \text{lkg}(v) \) together with
\( v \). Then the number of \((k+1)\)-cliques of \( G \) containing \( v \) is \( c_k(\text{lkg}(v)) \). In particular,
the choice of \( v_0 \) gives \( c_k(\text{lkg}(v_0)) \geq c_k(\text{lkg}(v')) \) for every vertex \( v' \in G \).

Define graphs \( G_0, G_1, \ldots, G_{s+1} \) by setting \( G_{i+1} = G - \{v_0, v_1, \ldots, v_i\} \) for \( 0 \leq i \leq s \)
and \( G_0 = G \). Clearly, \( G = G_0 \supset G_1 \supset \cdots \supset G_{s+1} \). Further, \( G_{s+1} \) is the induced
subgraph on the vertices adjacent to \( v_0 \), which is \( \text{lkg}(v_0) \).

Since \( c_{r+1}(G) = 0 \), \( c_r(\text{lkg}(v_0)) = 0 \), for otherwise, the \( r \) vertices of an \( r \)-clique
of \( \text{lkg}(v_0) \) together with \( v_0 \) would form an \((r+1)\)-clique of \( G \). Then \( c_r(G_{s+1}) = 0 \).

Further, since \( c_{k+1}(G) > 0 \), and \( v_0 \) is contained in the most \((k+1)\)-cliques of
any vertex of \( G \), \( v_0 \) is contained in at least one \((k+1)\)-clique, and so \( c_k(\text{lkg}(v_0)) > 0 \).

Since \( v \) is contained in at least one \((k+1)\)-clique of \( G \), we have \( c_k(\text{lkg}(G_{s+1})) < c_{k+1}(G) \).

Then by the second inductive hypothesis, there is some \((r-1)\)-colorable complex
\( C_{s+1} \) such that \( c_k(C_{s+1}) = c_k(G_{s+1}) \) and \( c_{k+1}(C_{s+1}) = c_{k+1}(G_{s+1}) \). Since given
any \((r-1)\)-colorable complex, there is an \((r-1)\)-colorable rev-lex complex with
the same face numbers, we can take \( C_{s+1} \) to be a rev-lex complex. Further, since
\( c_{k+1}(C_{s+1}) \) and \( c_k(C_{s+1}) \) only force a lower bound on \( c_{k-1}(C_{s+1}) \), but not an upper
bound, we can take \( c_{k-1}(C_{s+1}) \geq c_{k-1}(G) \).
Let \( c_k lk_G( v_i ) = a_i \) and \( c_k -1( lk_G( v_i ) ) = b_i \). Since \( G_{i+1} = G_i - v_i, c_{k+1}( G_i ) \) and \( c_k( G_i ) = c_k( G_{i+1} ) = a_i \) and \( c_k -1( G_i ) = c_k( G_{i+1} ) = b_i \). We have \( c_k( lk_G( v_i ) ) \geq c_k( lk_G( v_i ) ) \) by the choice of \( v_i \). We also have \( c_k( lk_G( v_i ) ) \geq c_k( lk_G( v_i ) ) \) since \( G_i \subset G \).

Thus
\[
c_k( C_{s+1} ) = c_k( G_{s+1} ) = c_k( lk_G( v_0 ) ) \geq c_k( lk_G( v_i ) ) \geq c_k( lk_G( v_i ) ) = a_i .
\]

Given an \( r \)-colored complex \( C_{i+1} \) such that \( c_{k+1}( C_{i+1} ) = c_{k+1}( G_{i+1} ) = c_k( G_{i+1} ) \) and the induced subcomplex of \( C_{i+1} \) on the vertices of the first \( r-1 \) colors is isomorphic to \( C_{s+1} \), we want to construct a complex \( C_i \) such that \( c_{k+1}( C_i ) = c_{k+1}( G_i ) \), and the induced subcomplex of \( C_i \) on the vertices of the first \( r-1 \) colors is isomorphic to \( C_{s+1} \).

Construct \( C_i \) from \( C_{i+1} \) by adding a new vertex \( v_i ' \) of color \( r \). Let the \( (k+1) \)-faces containing \( v_i ' \) consist of each of the first \( a_i \) \( k \)-faces in rev-lex order of \( C_{s+1} \) together with \( v_i ' \), and let the \( k \)-faces containing \( v_i ' \) consist of each of the first \( b_i \) \((k-1)\)-faces in rev-lex order of \( C_{s+1} \) together with \( v_i ' \).

If this construction can be done, then \( c_{k+1}( C_i ) \) is the number of \((k+1)\)-faces of \( C_i \) containing \( v_i ' \) plus the number of \((k+1)\)-faces of \( C_i \) not containing \( v_i ' \), which are \( a_i \) and \( c_{k+1}( C_{i+1} ) \), respectively. Then
\[
c_{k+1}( C_i ) = c_{k+1}( C_{i+1} ) + a_i = c_{k+1}( G_{i+1} ) + a_i = c_{k+1}( G_i ) .
\]

Likewise, we have
\[
c_k( C_i ) = c_k( C_{i+1} ) + b_i = c_k( G_{i+1} ) + b_i = c_k( G_i ) .
\]

Further, it is clear from the construction that the induced subcomplex on vertices of the first \( r-1 \) colors is unchanged from \( C_{i+1} \), and hence is isomorphic to \( C_{s+1} \).

In order to show that the construction is possible, we must show that \( c_k( C_{s+1} ) \geq a_i \) and \( c_k -1( C_{s+1} ) \geq b_i \), and that it is possible for an \((r-1)\)-colored complex \( C \) to have exactly \( c_k( C ) = a_i \) and \( c_k -1( C ) = b_i \). For the first of these, we have already shown that \( c_k( C_{s+1} ) \geq a_i \).

For the second, \( c_k -1( lk_G( v_i ) ) = b_i \). But \( lk_G( v_i ) \subset G_i \subset G \), so
\[
b_i = c_k -1( lk_G( v_i ) ) \leq c_k -1( G_i ) \leq c_k -1( G ) \leq c_k -1( C_{s+1} ) .
\]

For the third, since \( G_i \subset G \), we have \( c_{r+1}( G_i ) \leq c_{r+1}( G ) = 0 \), and so \( c_{r+1}( G_i ) = 0 \). Then \( c_r( lk_G( v_i ) ) = 0 \). We also have \( c_k( lk_G( v_i ) ) = a_i \) and \( c_k -1( lk_G( v_i ) ) = b_i \) by the definitions of \( a_i \) and \( b_i \). Then by the first inductive hypothesis, there is some \((r-1)\)-colored complex \( C_i ^\prime \) such that \( c_k( C_i ^\prime ) = a_i \) and \( c_k -1( C_i ^\prime ) = b_i \). Then we can take \( C_i ^\prime \) to be the \((r-1)\)-colored rev-lex complex with \( c_k( C_i ^\prime ) = a_i \) and \( c_k -1( C_i ^\prime ) = b_i \). Since \( C_{s+1} \) is an \((r-1)\)-colored rev-lex complex with \( c_k( C_{s+1} ) \geq a_i \) and \( c_k -1( C_{s+1} ) \geq b_i \), \( C_i ^\prime \subset C_{s+1} \), and we can choose the link of \( v_i ^\prime \) in \( C_i \) to be \( C_i ^\prime \).

We can repeat this construction for each \( 0 \leq i \leq s \) to start with \( C_{s+1} \), then construct \( C_s \), then \( C_{s-1} \), and so forth, until we have an \( r \)-colored complex \( C_0 \) such that \( c_k( C_0 ) = c_k( G ) \) and \( c_k -1( C_0 ) = c_k -1( G ) \). This completes the inductive step for the induction on \( c_k -1( G ) \), which in turn completes the inductive step for the induction on \( k \).

We are now ready to prove the result which immediately implies Theorem 1 and hence establish the Kalai-Eckhoff conjecture, by taking \( r \) to be the clique number of \( G \).

**Theorem 4.2.** For every graph \( G \) with \( c_{r+1}( G ) = 0 \), there is an \( r \)-colorable complex \( C \) such that \( c_i( C ) = c_i( G ) \) for all \( i \).
Proof. By Lemma 4.1, we can pick an $r$-colored complex $C_i$ such that $c_i(C_i) = c_i(G)$ and $c_{i+1}(C_i) = c_{i+1}(G)$ for all $i \geq 1$. By Theorem 3.7, we can take $C_i$ to be the rev-lex complex on $c_i(G)$ $i$-faces and $c_{i+1}(G)$ $(i+1)$-faces, and then $\bigcup_{i=1}^r C_i$ will have the desired face numbers. 

Acknowledgements. I would like to thank my thesis advisor Isabella Novik for suggesting this problem, and for her many useful discussions on solving the problem and writing an article.

References

[1] J. Eckhoff, Intersection properties of boxes. I. An upper-bound theorem, Israel J. Math. 62 (1988) no. 3 283-301.
[2] J. Eckhoff, The maximum number of triangles in a $K_4$-free graph, Discrete Math. 194 (1999) no. 1-3 95-106.
[3] J. Eckhoff, A new Turán-type theorem for cliques in graphs, Discrete Math. 282 (2004) 113-122.
[4] P. Frankl, Z. Füredi, and G. Kalai, Shadows of colored complexes, Math. Scand. 63 (1988) 169-178.
[5] G. Katona, A theorem of finite sets, in: Theory of Graphs, Academic Press, New York, 1968, pp. 187-207.
[6] J.B. Kruskal, The number of simplices in a complex, in: Mathematical Optimization Techniques, University fo California Press, Berkeley, California, 1963, pp. 251-278.
[7] R. Stanley, Balanced Cohen-Macaulay complexes, Trans. Amer. Math. Soc. 249 (1979) 139-157.
[8] R. Stanley, Combinatorics and Commutative Algebra, Second Edition, Birkhauser Boston, Inc., Boston, Massachusetts, 1996, 53-64.
[9] P. Turán, Eine Extremalaufgabe aus der Graphentheorie Mat. Fiz. Lapok 48 (1941) 436-452.
[10] A.A. Zykov, On some properties of linear complexes, Amer. Math. Soc. Transl. (1952) no. 79.

Department of Mathematics, University of Washington, Seattle, WA 98195-4350
E-mail address: frohmade@math.washington.edu