The Primal Framework I.

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This the first of a series of articles dealing with abstract classification theory. The apparatus to assign systems of cardinal invariants to models of a first order theory (or determine its impossibility) is developed in [5]. It is natural to try to extend this theory to classes of models which are described in other ways. Work on the classification theory for nonelementary classes [8] and for universal classes [9] led to the conclusion that an axiomatic approach provided the best setting for developing a theory of wider application. This approach is reminiscent of the early work of Fraissé and Jónson on the existence of homogeneous-universal models. As this will be a long project it seems appropriate to report our progress as we go along.

In large part this series of articles will parallel the development in [9]. A survey of that paper which could serve as an introduction to this one is [1]. The first chapter of this article corresponds to Section 2 of [9]. In it we describe the axioms on which the remainder of the article depends and give some examples and context to justify this level of generality. As is detailed later the principal goal of this series is indicated by its title. The study of universal classes takes as a primitive the notion of closing a subset under functions to obtain a model. We replace that concept by the notion of a prime model. We begin the detailed discussion of this idea in Chapter II. One of the important contributions of classification theory is the recognition that large models can often be analyzed by means of a family of small models indexed by a tree of height at most $\omega$. More precisely, the analyzed model is prime over such a tree. Chapter III provides sufficient conditions for prime models over such trees to exist. The discussion of properties of a class which guarantee that each model in the class is prime over such a tree will appear later in the series.

We introduce in Chapters I and II a number of principles, which we loosely refer to as axioms. At the beginning of Chapter III we define the notion of an adequate class — a class which satisfies those axioms that we assume in the mainline of the study. This notion of an adequate class will be embellished by further axioms in later papers of this series. Our use of the word axiom in this context is somewhat inexact; postulate might be better. In exploring an unknown area we list certain principles which appear to make important distinctions. In our definition of an adequate class we collect a family of these principles.
that is sufficient establish a coherent collection of results. We thank N. Shi for carefully reading the paper and making a number of helpful suggestions.
Chapter I

THE ABSTRACT FRAMEWORK

Shelah developed in [9] several frameworks for studying aspects of classification theory. In each case he studied a triple

\[ K = \langle K, \leq_K, NF_K \rangle; \]

where \( K \) is a collection of structures, \( \leq_K \) denotes elementary submodel with respect to \( K \), and \( NF_K \) is a 4-ary relation (nonforking) denoting that certain models are in stable amalgamation. The original paper primarily studied classes which admitted a fourth basic notion: ‘generated submodel’. We generalize that context here by taking as a fourth basic component a predicate \( cpr \). Intuitively, \( cpr = cpr_K \) holds of a structure \( M \in K \) and a chain of models \( M \) (Section I.2) if \( M \) is prime over \( M \).

Thus this paper studies quadruples

\[ K = \langle K, \leq_K, NF_K, cpr_K \rangle. \]

Section 1 reviews the properties of elementary submodel and free amalgamation which carry over from [9]. In section 2 we provide a number of examples of classes which satisfy the basic axioms.

I.1 Basic properties of \( \leq_K \) and \( NF \)

\( K \) always denotes a class of structures of a fixed similarity type. \( K \) and all relations that we define on it are assumed to be closed under iso-
morphism. Although technically both $\leq$ and $\text{NF}$ should be subscripted with $K$, we usually omit the subscript for ease of reading. $M$ and $N$ (with subscripts) denote members of $K$ unless we explicitly say otherwise. $A$ and $B$ will denote subsets of members of $K$. We write $M$ is contained in $N$ ($M \subseteq N$) if $M$ is a substructure of $N$; that is, if the universe of $M$ is a subset of that of $N$, the relations of $N$ are those imposed by $M$ and $M$ is closed under any functions in the language. We write $M$ is a submodel or more explicitly a $K$-submodel of $N$ for $M \leq N$.

The first group of axioms describe our notion of elementary submodel.

I.1.1 Axiom Group A: $K$-submodels.

A0 If $M \in K$ then $M \leq M$.

A1 If $M \leq N$ then $M$ is a substructure of $N$.

A2 $\leq$ is transitive.

A3 If $M_0 \subseteq M_1 \subseteq N$, $M_0 \leq N$ and $M_1 \leq N$ then $M_0 \leq M_1$.

It is sometimes essential to distinguish between $M \leq K N$ which implies $M \subseteq N$ and the existence of an embedding $f$ of $M$ into $N$ whose image is a $K$-submodel of $N$.

I.1.2 Definition. A $K$-embedding is an isomorphism $f$ from an $M$ in $K$ to an $N$ in $K$ such that $\text{rng} f \leq K N$. $N$ is then a $K$-extension of $M$ via $f$.

If $f$ is not mentioned explicitly then it is the identity.

We require one further important property of $K$ and $\leq K$.

I.1.3 Definition. $K$ has the $\lambda$-Löwenheim-Skolem property ($\lambda$-LSP) or (the Löwenheim-Skolem property down to $\lambda$) if $A \subseteq M \in K$ and $|A| \leq \lambda$ implies there is an $N \in K$ with $A \subseteq N \leq M$ and $|N| \leq \lambda$. The Löwenheim-Skolem number of $K$ is the least $\lambda$ such that $K$ has the $\lambda$-Löwenheim-Skolem property. We write $\text{LS}(K) = \lambda$.

Note that the set of cardinals for which $K$ has the Löwenheim Skolem property may not be convex. Moreover the related requirement on $\lambda$, for any $A$ there is an $N \supseteq A$ with $N \in K$ and $|N| \leq |A| + \lambda$, is still different and will be investigated later.
I.1. BASIC PROPERTIES OF $\leq_K$ AND NF

I.1.4 Axiom Group A: K-submodels.

A4 $LS(K) < \infty$.

Our axioms differ from those of Fraissé and Jónsson in that there is neither a joint embedding nor an amalgamation axiom. Our approach here is to assume in the next set of axioms a particularly strong form of amalgamation. The use of nonamalgamation as a source of nonstructure has been explored by Shelah in several places. See especially Chapter I of [3] and [4] and its progenitor [8]. We will obtain joint embedding by fiat (by restricting to a subclass that satisfies it). The Jónsson Fraissé constructions also require closure under unions of chains. This requirement is more subtle than it first appears; it is the major topic of [3].

We say two structures are compatible if they (isomorphic copies of them) have a common $K$-extension. Any class $K$ that satisfies Axiom C2 (below) is split into classes with the joint embedding property by the equivalence relation of compatibility and if $K$ has a Löwenheim-Skolem number then there will be only a set of equivalence classes. Given any fixed model $M$ (or diagram of models) this equivalence relation will be refined by 'compatibility over $M$'. We explore this refinement in Chapter II.

The second group of axioms concern the independence relation. We begin by describing a relation of four members of $K$.

$$NF(M_0, M_1, M_2, M_3)$$

means that $M_1$ and $M_2$ are freely amalgamated over $M_0$ within $M_3$. The notation NF arises from the reading the type of $M_1$ over $M_2$ inside $M_3$ does not fork over $M_1$. We will eventually show a dichotomy between nonstructure results and the existence of a monster model $\mathcal{M}$. Thus, in trying to establish a structure theory we can introduce a 3-ary relation $NF(M_0, M_1, M_2)$ to abbreviate $NF(M_0, M_1, M_2, \mathcal{M})$. We usually write this relation as $M_1 \downarrow_{M_0} M_2$. Even before showing the existence of the monster model we will employ this notation if the choice of $M_3$ is either clear from context or there are several possibilities which serve equally well.

I.1.5 Notation. A 4-tuple $\langle M_0, M_1, M_2, M_3 \rangle$ is called a full free amalgam if it satisfies $NF(M_0, M_1, M_2, M_3)$. The three-tuple $\langle M_0, M_1, M_2 \rangle$
is called a free amalgam if it is an initial segment of a full free amalgam. We often write $M_1$ and $M_2$ are freely amalgamated over $M_0$ in $M_3$. We refer to such a diagram as a ‘free vee’. An isomorphism between two free amalgams $M = \langle M_0, M_1, M_2 \rangle$ and $M' = \langle M'_0, M'_1, M'_2 \rangle$ is a triple of isomorphisms $f_i$ mapping $M_i$ to $M'_i$ with $f_0$ contained in $f_1$ and $f_2$. There is no guarantee (until we assume Axiom C5 below) that the isomorphisms $f_1$ and $f_2$ have a common extension to an $M_3$ which completes $M$. We extend this notion of isomorphism to arbitrary diagrams in Section 2.

I.1.6 Axiom Group C: Independence. The following axioms describe the independence relation. For convenience of comparison we have kept the numbering from [9] when we have just copied an axiom. Some axioms from that list (e.g. C4) are omitted here. In particular, the role of Axiom Group B from [9], which dealt with the notion of generation, is played here by Axiom Group D. (See Section I.2.)

C1 If $NF(M_0, M_1, M_2, M_3)$ then $M_0 \leq M_2 \leq M_3$ and $M_0 \leq M_1 \leq M_3$. In particular, each $M_i \in K$. If $NF(M_0, M_1, M_2, M_3)$, we say $M_1$ and $M_2$ are freely amalgamated (or independent) over $M_0$ in $M_3$.

C2 Existence If $M_0$ is a $K$-submodel of both $M_1$ and $M_2$ then there are copies (over $M_0$) $M'_1$ and $M'_2$ of $M_1$ and $M_2$ which are freely amalgamated in some $M_3 \in K$.

C3 Monotonicity i) If $M_1$ and $M_2$ are freely amalgamated over $M_0$ in $M_3$ then so are $M_1$ and $M'_2$ for any $M'_2$ with $M_0 \leq M'_2 \leq M_2$.

ii) If $M_1$ and $M_2$ are freely amalgamated over $M_0$ in $M_3$ then they are freely amalgamated in any $M'_3 \geq M_3$.

iii) If $M_1$ and $M_2$ are freely amalgamated over $M_0$ in $M_3$ then they are freely amalgamated in any $M'_3$ containing $M_1 \cup M_2$ and with $M'_3 \leq M_3$.

C5 Weak Uniqueness Suppose $\langle M, M_3 \rangle$ and $\langle M', M'_3 \rangle$ are full free amalgams. If $M$ and $M'$ are isomorphic free amalgams (via $f$) then there is an $N \in K$ with $M'_3 \leq N$, and an extension of $f$ mapping $M_3$ isomorphically onto a $K$-submodel of $N$. 
I.1. BASIC PROPERTIES OF \( \leq_k \) AND NF

C6 Symmetry If NF\((M_0, M_1, M_2, M_3)\) then NF\((M_0, M_2, M_1, M_3)\).

C7 Disjointness If NF\((M_0, M_1, M_2, M_3)\) then \( M_1 \cap M_2 = M_0 \).

Axiom C7 is largely a matter of notational convenience. We will indicate in Chapter III how the major argument of this paper could be slightly revised to avoid this axiom. With it we obtain immediately the following monotonicity property.

I.1.7 Proposition. Suppose NF\((M_0, M_1, M_2, M_3)\) and \( M_0 \leq N \) which is a \( K \)-submodel of \( M_1 \) and \( M_2 \). Then NF\((N, M_1, M_2, M_3)\).

This result could be easily deduced from the base extension axiom II.2.2 without using C7.

By taking ‘formal copies’ of \( M'_2 \) and \( M_3 \), one derives a variant of C2 where \( M'_1 \) can be demanded to be \( M_1 \). We refer to Axiom C5 as weak uniqueness because it simply demands that any two amalgamations of a given vee be compatible. Thus, it is making the compatibility class of the diagram unique, not the amalgamating model.

I.1.8 Lemma. Suppose \( \langle M_0, M_1, M_2, M_3 \rangle \) and \( \langle M_0, M_1, M'_2, M'_3 \rangle \) are full free amalgams. If \( M'_2 \) is isomorphic to \( M_2 \) over \( M_0 \) by a map \( g \) then \( g \) is an isomorphism between \( M'_2 \) and \( M_2 \) over \( M_1 \).

Proof. Apply the weak uniqueness axiom to the map \( f \) that is the union of the identity on \( M_1 \) and the given \( g \) from \( M_2 \) to \( M'_2 \).

I.1.9 Smoothness. Does the class \( K \) admit a ‘limit’ of an ascending chain (or more generally a directed system) of \( K \)-stuctures? There are several different variants on this question and the answers determine significant differences in the behavior of \( K \). We discuss the variants in detail in [3]; we now just mention a couple of possibilities and some of the consequences.

The strongest requirement is to deal directly with unions of chains. But even here there are several variations. One can demand that any union of a (continuous) increasing chain of \( K \) stuctures be a member of \( K \). More subtly, one can ask that if each member of the chain is \( K \)-submodel of a fixed \( M \) then the union is also.
Shelah has shown that a class $K$ satisfying the axioms A0 through A4 enumerated here and stringent requirements for closure under unions can be presented as the collection of models in a pseudoelementary class in an infinitary logic which omit a family of types. (See Section 1 of [8] and [4].)

Beginning in Chapter II we discuss the ways in which unions of chains can be replaced by demanding the existence of a prime model over the union. Again, there are a number of smoothness properties that can be discussed in this context. One obvious application is attempts to improve the Löwenheim Skolem property to demand that each set can be imbedded in a model of roughly the same size.

I.2 Examples

We describe in this section a number of examples of classes and notions of amalgamation which satisfy at least some of the axioms that we are discussing. Of course, the prototype is the collection of models of a stable first order theory where a free amalgamation is one that is independent in the sense of nonforking.

In this section we first discuss some contrived examples which although they lack any intrinsic interest make it easy to exhibit some of the pathologies that we are investigating. Then we place in context some classes which naturally arise in the attempt to extend classification theory to, e.g., infinitary classes.

I.2.1 Contrived Examples. Let $B$ be the class of all structures of the following sort. We work in a language with two unary predicates, $U$, $V$ and a binary relation $E$. Now $M$ is in $B$ if via $E$, each member $a$ of $V$ is the name of a subset $X_a = \{m \in U(M) : E(m, a)\}$ of $U$ and every subset of $U$ has one and only one name. Thus each member $M$ of $B$ is determined up to isomorphism by the cardinality of $U(M)$. $NF(M_0, M_1, M_2, M_3)$ holds just if for each $a \in M_1 \ (M_2)$, $X_a$ in the sense of $M_3$ is a subset of $M_1 \ (M_2)$. We can illustrate the axioms by defining $\leq_B$ in two different ways.

i) Define $M \leq_B N$ if $U(M) \subseteq U(N)$, $V(M) \subseteq V(N)$ and each element of $V(M)$ names in $N$ the same subset of $U(M)$ that
it names in $M$. Under this definition if $M_0 \leq_B M_1 \leq_B M_3$ and $M_0 \leq_B M_2 \leq_B M_3$ then $\text{NF}(M_0, M_1, M_2, M_3)$. Moreover if $M_0 \leq_B M_1$ and $M_0 \leq_B M_2$, we can find a common $K$-extension for them letting $M_3$ be $M_1 \cup M_2$ together with a collection of names for sets that intersect both $U(M_1)$ and $U(M_2)$.

ii) On the other hand, let $M \leq_B N$ just mean that $M$ is a substructure of $N$. Now it is still possible to verify Axiom C2. If a subset $X$ of $M_0$ is named by elements $a$ of $M_1$ and $b$ of $M_2$ then $a$ and $b$ can be identified by the embeddings into $M_3$. Using this strategy to amalgamate Axiom C7 fails; however the strategy outlined below when we consider an additional predicate $Q$, which is needed then to obtain even C2, will also show that C7 is verified.

Some of the problems with amalgamations become clear if we add a unary predicate $Q$ and demand that for any subset $W$ of $U(M)$ with power less than $\kappa$ there exist both $a$ and $b$ in $V(M)$ with $a$ satisfying $Q$ but $b$ not satisfying $Q$ so that $X_a$ and $X_b$ both contain $W$. (That is, both $Q$ and its complement are dense.) Axioms C2 and C5 hold under the first definition of $\leq_B$. To make C2 hold under the second interpretation of $\leq$ we must deal with a subset $X$ of $M_0$ that has one name in $Q(M_1)$ and another in $\neg Q(M_2)$. Now another strategy works. Add points to the set attached to one of the names and then fill out the model as freely as possible. It is easy to see (cf. Chapter II) that there are extensions of chains which are incompatible.

$B_\kappa$ is defined in the same way but with the additional restriction that each $|X_a| < \kappa$.

Another variant arises by replacing the single binary relation $E$ by a family of binary relations $E_i$ such that for each $i < \kappa$ and each $a \in V(M)$, there is a unique $m \in U(M)$ with $R_i(a, m)$. Thus we code $\kappa$ sequences rather than sets.

**I.2.2 $\aleph_1$-saturated models.** Let $T$ be a strictly stable first order theory and let $K$ be the class of $\aleph_1$-saturated models of $T$. Take $\text{NF}_K$ as nonforking and $\leq_K$ as elementary submodel. Then the basic axioms are clearly satisfied but $K$ is not closed under unions of chains of small cofinality (not fully smooth). But there are prime models ($\mathbf{F}_\aleph_1$ in the notation of [3] or $\mathbf{SET}_{\aleph_1}$ in the notation of [4]) over such chains.
The theory $\text{REF}_\omega$ of countably many refining equivalence relations has $K$-prime models of over chains of cofinality $\omega$ but they are not minimal. This leads to $2^\lambda$ $K$-models of power $\lambda$ when $\lambda^\omega = \lambda$. This argument is treated briefly in [6] (the didip) and will be reported at more length in the current series of papers.

On the other hand if $T$ is a two dimensional stable theory (cf. Theorem V.5.8 of [5]) then $I(\aleph_\alpha, K) \leq |\alpha + 1|$.

I.2.3 Universal Classes. See II.2.2 of [5].

I.2.4 Finite diagrams stable in power. See II.2.3 of [5].

I.2.5 Infinitary Classes. See [8].

I.2.6 Banach Spaces. See [7].

1 Question. Are there are stable universal theories of Banach Spaces beyond the $L^p$-spaces?
Chapter II

Prime models over diagrams

In Section 1 we discuss diagrams of models and the basic properties of prime models over diagrams. Section 2 concerns prime models over independent pairs of models. We also discuss several possibilities for the relation of a model $M_1$ which is independent from an $M_4$ over $M_0$ with an $M_2$ with $M_0 \subseteq M_2 \subseteq M_4$. Surprisingly, II.2.3 deduces a property of the dependence relation but needs the properties of prime models for the proof. In Section 3 we discuss prime models over chains.

II.1 Diagrams

We consider two kinds of diagrams. The first is more abstract because the partial ordering among the structures is witnessed by $K$-embeddings. In the (second) case of a concrete diagram the partial ordering is witnessed by actual set theoretic containments. The existence axioms for free amalgamations have the general form: Given an abstract diagram, there is a concrete stable diagram which is isomorphic to it (in the category of diagrams).

The discussion of abstract diagrams is essential at this stage in the development of theory. Once a ‘monster’ model or global universe of discourse has been posited, one can assume that all diagrams are concretely realized as subsets of the monster model. The monster model is easily justified in the first order context so this subtlety does not seriously arise. In the more general case one must worry about it until
‘smoothness’ and thus the monster model are obtained. (See § 3 and 3).

We will assume the existence of prime models over certain simple diagrams and certain properties relating prime models and independence. From this we will deduce the existence of prime models over more complicated diagrams. In a later paper when considering prime model over still more complicated diagrams (tall trees) we will only be able to obtain a dichotomy between the existence of prime models and the existence of many models.

We call a triple \( \langle M_0, M_1, M_2 \rangle \) with embeddings of \( M_0 \) into \( M_1 \) and \( M_2 \) a ‘vee’ diagram. Thus, Axiom C2 asserts the existence of free amalgamations over vee’s.

II.1.1 Definition. An abstract \( K \)-diagram (indexed by a partial order \( \langle I, \preceq \rangle \)) is a pair: a sequence of models and a collection of maps. The models \( \{ M_x : x \in I \} \) and the maps \( \{ f_{xy} : x, y \in I, x \preceq y \} \) must satisfy the following conditions.

- \( M_x \in K \).
- If \( x \preceq y \) then there is a \( K \)-embedding \( f_{xy} \) from \( M_x \) into \( M_y \).
- If \( x \preceq y \preceq z \) then \( f_{xz} = f_{yz} \circ f_{xy} \).

We will often suppress mention of the maps. When we need to refer to them we will often describe the family \( \{ f_{xy} : x, y \in I \} \) by \( f \) and write \( f|Y \) for \( \{ f_{xy} : x, y \in Y \} \) when \( Y \) is a subset of \( I \).

II.1.2 Definition. \( \underline{M} \) and \( \underline{N} \) are isomorphic \( K \)-diagrams if they are both indexed by the same partial order \( I \) and for each \( x \in I \) there is an isomorphism \( \alpha_x \) between \( M_x \) and \( N_x \) such that the \( \alpha_x \) and \( f_{xy} \) in \( \underline{M} \) and \( \underline{N} \) commute in the natural way. We will write \( \alpha \) for the family of the \( \alpha_x \).

Note that no single function \( f \) is a morphism for \( K \) unless its domain is in \( K \). In particular, we cannot speak of a single map whose domain is a vee.

II.1.3 Definition. A concrete \( K \)-diagram \( \underline{M} \) inside \( M \) (indexed by \( I \)) is a collection, \( \underline{M} = \{ M_x : x \in I \} \), of members of \( K \) such that each \( M_x \leq M \) and if \( x \preceq y \) then \( M_x \leq M_y \).
We extend the notion of a free amalgam to more general diagrams.

**II.1.4 Definition.** A *stable* $K$-diagram $M$ inside $M$ (indexed by the lower semilattice $\langle I, \land \rangle$) is a *concrete* $K$-diagram inside $M$ satisfying the additional condition

- For any $x, y \in I$, $M_x$ and $M_y$ are freely amalgamated over $M_{x \land y}$ within $M$.

**II.1.5 Definition.**

i) $N$ is a $K$-*extension* of $M$ if there is an isomorphism $f$ between $M$ and a concrete $K$-diagram $N$ inside $N$. $f$ is called a $K$-*embedding* of $M$ into $N$.

ii) $N$ is a *stable* $K$-*extension* of $M$ if there is an isomorphism $f$ between $M$ and a stable $K$-diagram $N$ inside $N$. $f$ is called a *stable* $K$-*embedding* of $M$ into $N$.

The phrase ‘$M$ is a stable $K$-diagram’ means $M$ is stable in some $M$ which is suppressed for convenience. A vee $\langle M_0, M_1, M_2 \rangle$ is stable if there is an $M_3$ completing it to a full free amalgam. Note that an isomorphism between diagrams need not preserve stability. It only guarantees isomorphism between the individual $M_x$ and $M_x'$. In particular an isomorphic copy of a free amalgam need not be free. So we cannot say a vee diagram is free without specifying an ambient model and an embedding into it. This is true even in the first order case. In this case we usually think of the monster model as the ambient model and identity as the embedding.

**II.1.6 Examples.** Of course a (full) free amalgam is a stable $K$-diagram. So is any chain of $K$-models which have a common $K$-extension. Another natural example is an ‘L’-diagram (see figure) where for each $i \leq \delta$, $M_i \downarrow_{M_0} N$ inside the ambient model.

Compatibility of two extensions of a single model is a straightforward notion.

**II.1.7 Definition.** Two members $N_1, N_2$ of $K$ are said to be *compatible* over the $K$-embeddings $f_1, f_2$ of $N_0$ into $N_1, N_2$ if there is an $N_3$ in $K$
Figure II.1: An ‘L’ diagram
with both $N_1$ and $N_2$ $K$-embeddible in $N_3$ by $K$-embeddings $g_1, g_2$ such that $g_1 \circ f_1 = g_2 \circ f_2$.

In view of the existence axiom for free amalgamations any two members $M_1, M_2$ of $K$ with an $M_0 \in K$ that is a $K$-substructure of each of them are compatible over $M_0$. When $N_0$ is replaced by a diagram $M$ involving infinitely many members of $K$ the situation is more complicated.

II.1.8 Definition. Suppose $M$ is a $K$-diagram. Let $f_1, f_2$ be $K$-embeddings of $M$ into $M_1$ and $M_2$ respectively. $M_1$ and $M_2$ are $K$-compatible over $M$ via $f_1, f_2$ if there exists an $N \in K$ and $K$-embeddings $g_1, g_2$ of $M_1, M_2$ into $N$ such that $g_1 \circ f_1$ and $g_2 \circ f_2$ agree on $M$.

(The $f_i$ are families of maps and the $g_i$ are single maps but the meaning of the composition should be clear.)

Clearly over each $M$, compatibility defines a reflexive and symmetric relation. Axiom C2 shows that the relation is also transitive so for any diagram $M$ we have an equivalence relation, compatibility over $M$.

II.1.9 Definition. The abstract $K$-diagram $M$ is stably univalent if all stable $K$-extensions of $M$ are compatible.

In this language Axiom C5 asserts that every vee is stably univalent.

There is little hope to find a ‘prime’ model (in the usual categorical sense) over an arbitrary diagram $M$. For, if a diagram isomorphism from $M$ into $N$ collapsed two elements of models $M \in M$ then there could be no isomorphism from a structure containing all members of $M$ into $N$. Thus, we restrict the following definitions to stable embeddings. We first define the notion of a ‘prime’ model within a compatibility class. Remember that the image of a stable embedding is required to be stable.

II.1.10 Definition. i) Let $f$ be a stable $K$-embedding of the abstract diagram $M$ in $M$. We say $M$ is compatibility prime over $(M, f)$ (or over $M$ via $f$) if for every $M' \in K$ and stable embedding $f'$ of $M$ into $M'$ such that $M$ and $M'$ are compatible via $f, f'$ there is an embedding of $g : M \mapsto M'$ with $g \circ f = f'$.

ii) We omit $f$ if it is an identity map.
Strictly speaking, we should refer to a triple \((\mathcal{M}, f, M)\). A little looseness to ease reading seems acceptable here. We are more precise when we introduce the notion of canonically prime in Section II.3 since a new basic relation is added to the system.

Note that if the diagram \(\mathcal{M}\) has a unique maximum element then that element is compatibility prime over \(M\). Note that if \(M\) is compatibility prime over \(\mathcal{M}\) via \(f\) then \(M\) is isomorphic to an \(M'\) which is compatibility prime over \(\mathcal{M}\) via the identity.

There are various ways in which a compatibility prime model can fail to be unique. There could be more than one compatibility class; within a given compatibility class there could be nonisomorphic compatibility prime models.

II.1.11 Definition. \(M\) is absolutely prime over \((\mathcal{M}, f)\) if \(M\) can be embedded over \((\mathcal{M}, f)\) into any stable \(K\)-extension \(N\) of \(M\).

Clearly if \(M\) is compatibility prime over \(\mathcal{M}\) and \(M\) is univalent then \(M\) is absolutely prime. In view of the weak uniqueness axiom if there is a prime model over a stable diagram \(\mathcal{M}\) then \(M\) is stably univalent. We will see that with strong enough hypotheses on \(K\) the various notions of prime coalesce. We have introduced the notion of ‘absolutely prime’ to emphasize the distinction with compatibility prime; ‘absolutely prime’ is the natural extension of the usual model theoretic notion to ‘prime over a diagram’. When we write ‘prime’ with no adjective we mean ‘absolutely prime’.

II.2 Prime Models over vees

In the next two sections we discuss axioms concerning the existence and properties of prime models over certain specific diagrams. We begin by assuming the existence of prime models over vee’s. In the light of C5 (weak uniqueness) all completions of an amalgam are compatible so it would be equivalent to replace absolutely prime by compatibility prime in the following axiom. We will often just say ‘prime’ for the absolutely prime model over a vee. Again because of weak uniqueness we don’t have to specify \(M_3\) is prime over \(M_1 \cup M_2\) in \(M_4\).

II.2.1 Prime Models over vees.
II.2. PRIME MODELS OVER VEES

Figure II.2: Base Extension Axioms

D1 There is an absolutely prime model over any free amalgam \(\langle M_0, M_1, M_2 \rangle\).

For compactness we write \(M\) is prime over \(M_1 \cup M_2\) instead of \(M\) is prime over \(\langle M_0, M_1, M_2 \rangle\).

Now we describe properties which relate independence and prime models. The next axiom of this group corresponds to Axiom C4 of [9]. Diagram II.2 illustrates each of the principles discussed below.

II.2.2 The base extension axiom.

D2 Suppose \(M_1\) and \(M_4\) are freely amalgamated over \(M_0\) in \(M_5\). If \(M_0 \leq M_2 \leq M_4\) and \(M_3\) is prime over \(M_1 \cup M_2\) in \(M_5\) then \(M_3\) and \(M_4\) are freely amalgamated over \(M_2\) in \(M_5\).

To understand the base extension axiom in the first order context, think of \(M_3\) as \(M_2[M_1]\) (For this notation see [2].). Then the axiom is implied by the fact that for any \(X\) and any model \(M\) (in the appropriate category, e.g. \(\omega\)-stable and the normal notion of prime model), \(X\) dominates \(M[X]\) over \(M\).

Axiom II.2.2 yields a somewhat surprising consequence. We can obtain the following ‘transitivity of nonforking’ from II.2.2 and the weak uniqueness we posited in Axiom C5. This result is remarkable because we are establishing a property of the nonforking relation which makes no reference to prime or generated models. But, the proof uses properties of either generated models (the version in [3]) or prime models (here).
**Theorem II.2.3 (Transitivity of independence).** If $\text{NF}(M_0, M_1, M_2, M_3)$ and $\text{NF}(M_2, M_3, M_4, M_5)$ then $\text{NF}(M_0, M_1, M_4, M_5)$.

Proof. By the existence axiom there are $M_4''$ and $M_5''$ and an isomorphism $g$ of $M_4$ and $M_5''$ over $M_0$ such that $M_1 \downarrow M_0 M_4''$ in $M_5''$. By monotonicity (C3 iii) we may assume $M_5''$ is prime over $M_1 \cup M_4''$. Let $M_5''$ denote $g(M_2)$ so $M_0 \leq M_5'' \leq M_4''$. We want to show the existence of an isomorphism with domain $M_5''$ which fixes $M_1$ and maps $M_4''$ to $M_4$.

Let $M_4'$ be prime over $M_1 \cup M_2$ and contained in $M_3$. By the monotonicity axioms (C3 iii) and C3 i) we have $M_1 \downarrow M_0 M_2$ in $M_3'$ and $M_3 \downarrow M_2 M_4$ in $M_5$. Similarly, if $M_5'$ is chosen prime over $M_3' \cup M_4$ in $M_5$ then $M_5' \leq M_5$ and $M_3' \downarrow M_4 M_3'$ in $M_5'$ using axiom C3 iii).

Since $M_2$ and $M_5''$ are isomorphic over $M_0$ and both are independent from $M_1$ over $M_0$ there is an isomorphism $f$ taking $M_3'$ into $M_5''$ which extends $g|M_2 \cup 1_{M_1}$. (This follows because $M_3'$ is prime over $M_2 \cup M_1$ and applying the weak uniqueness axiom.) Let $M_5''$ denote $f(M_3')$. Then $M_5''$ is prime over $M_2'' \cup M_1$ since $f$ is an isomorphism. By the base extension axiom II.2.2 $M_2'' \downarrow M_2 M_5''$ in $M_5''$. Let $M_5''$ be prime over $M_2'' \cup M_3$ and let $h$ be a map from $M_5''$ into some $N$ which extends $g \cup f^{-1}$. Now $M_4' = h(M_3'')$ and $M_4$ are isomorphic over $M_2$ and both are independent from $M_3'$ over $M_2$. So by the weak uniqueness axiom there is an $h'$ defined on $N$ which takes $M_3'$ to $M_4$ and fixes $M_3$. Now $h' \circ h$ takes $M_3''$ to $M_4$ and fixes $M_1$ as required.

**Theorem II.2.4 (Transitivity of primality).** Suppose $M_1$ and $M_4$ are freely amalgamated over $M_0$ in $M_5$. If $M_0 \leq M_2 \leq M_4$, $M_3$ is prime over $M_1 \cup M_2$ in $M_5$ and $M_5$ is prime over $M_3 \cup M_4$ in $M_5$ then $M_5$ is prime over $M_1 \cup M_4$ in $M_5$.

Proof. Let $N$ be a stable K-extension of $\langle M_0, M_1, M_4 \rangle$. That is, there exist maps $f_1$, $f_4$ from $M_1$, $M_4$ to K-submodels $M_1', M_4'$ of $N$ that agree on $M_0$; let $M_0'$ denote the common image of the $f_i$ on $M_0$. Then since the embedding is stable $M_1 \downarrow M_0' M_4'$ inside $N$. We must find a common extension of the $f_i$ to $M_5$. Let $f_2$ denote the restriction of $f_4$ to $M_2$ and $M_2'$ its image. Now, since $M_3$ is prime over $M_1 \cup M_2$ there is a map $f_3$ with domain $M_3$ which extends $f_1 \cup f_2$. Denote the image of $f_3$ by
\[ M_3' \]. By the base extension axiom \textit{II.2.2} \[ M_3' \downarrow M'_2 \] \[ M'_4 \] in \( N \). So there is an \( f_5 \) mapping \( M_5 \) into \( N \) and extending \( f_3 \cup f_4 \). A fortiori, \( f_5 \) extends \( f_3 \cup f_4 \) and we finish.

There are some further properties of prime models which both arise in some natural situations and are useful tools. We describe them now but they do not play an important role in the theory until we reach some rather special cases.

\textbf{II.2.5 Definition.} The concrete \( K \)-extension \( M \) of \( \bar{M} \) is \textit{minimal} over the diagram \( \bar{M} \) if there is no proper \( K \)-submodel of \( M \) that contains \( \bar{M} \).

\textbf{II.2.6 Some further Axioms.}

\textbf{D3} The prime model over a free amalgam \( M \) is minimal over \( M \).

\textbf{D4} The prime model over a free amalgam \( M \) is unique up to isomorphism over \( M \).

\textbf{D5} Suppose \( M_1 \) and \( M_4 \) are freely amalgamated over \( M_0 \) in \( M_5 \) and \( M_5 \) is prime over \( M_1 \cup M_4 \). If \( M_0 \preceq M_2 \preceq M_4 \) and \( M_3 \) is prime over \( M_2 \cup M_1 \) then \( M_5 \) is prime over \( M_4 \cup M_3 \).

D5 is a kind of ‘converse’ to Theorem \textit{II.2.4} that we may need later. It will only hold in very restricted cases. It is true for the notion of generation in universal classes; it fails for prime models in the first order case (even \( \omega \)-stable). However, if prime models over vees are minimal (e.g. for a first order theory without the dimensional order property ) then an even stronger version of D5 holds easily. Namely the monotonicity requirement that if \( M_5 \) is minimal over \( M_1 \cup M_4 \) then \( M_5 \) is minimal over \( M_3 \cup M_4 \).

\textbf{II.3 Prime models over chains}

An abstract chain \( \langle M_i, f_{i,j} : i, j \in I \rangle \) is an abstract diagram whose index set \( I \) is linearly ordered. Any closed initial segment of an abstract chain has a natural representation as a concrete chain. To see this consider
Let $M_i$ denote $f_{i,\alpha}(M_i)$. Then the $M_i'$ are a concrete chain with inclusion maps $f'_{i,j} = f_{j,\alpha} \circ f_{i,j} \circ f_{i,\alpha}^{-1}$. Thus for any chain indexed by an ordinal $\gamma$ and any limit ordinal $\delta < \gamma$ it makes sense to speak of $\cup_{i<\delta} M_i$ as we can concretely realize $M|\delta$ in $M_\beta$ for any $\beta$ with $\delta \leq \beta \leq \gamma$.

We discuss in this section the specification and existence of prime models over chains. In general, given an increasing chain of $K$-models there is no reason to assume that the chain has any common extension in $K$, let alone one that is prime. If we assume the existence but not the weak uniqueness of compatibility prime models, there may be two incompatible compatibility prime models over the same chain. (We say $K$ is not smooth.) For this reason, we can not in the most general case just define ‘prime’ as compatibility prime and posit that ‘prime’ models exist. We would need to introduce another axiom asserting that there is only one compatibility class over any chain. Justification of such an axiom is the main point of [3]. However, to reduce the set theoretic hypotheses of that argument we introduce a new predicate (cpr) with the intuitive meaning, ‘$M$ is canonically prime over $M'$ and prescribe axioms describing the behavior of such prime models.

Consider the last example in Paragraph I.2.1. We have a unary predicate $U$ which is dense and codense. At any limit stage of cofinality $\kappa$, we have to decide whether the name of certain $\kappa$ sequences are in $U$ or not. Different answers correspond to different compatibility classes.

We want to demand the existence of a ‘prime’ model over a union of a chain. If the chain has length longer than $\omega$ several possibilities arise for what we should demand of models at limit stages in the chain. It is unreasonable to demand the existence of ‘prime’ models over chains that are not $K$-continuous in the following sense.

**II.3.1 Definition.**

i) The chain $\langle M_i, f_{i,j} : i, j < \beta \rangle$ is $K$-continuous if for each limit ordinal $\delta < \beta$, cpr($M_\delta, M_\delta, f|\delta$).

ii) The chain $\langle M_i, f_i : i < \beta \rangle$ is continuous if for each limit ordinal $\delta < \beta$, $M_\delta = \cup_{i<\delta} M_i$.

Axiom Ch1 gives an implicit definition of a canonically prime model; Axiom Ch2 asserts that such a model exists. Suppose $M$ is a ‘long’ chain. At each limit stage one has a choice of compatibility classes.
II.3. PRIME MODELS OVER CHAINS

for a ‘prime’ model over that segment of the chain. The notion of a
canonically prime model requires that these choices cohere.

**Ch1** $\text{cpr}(\mathcal{M}, M, f)$ implies

1. $\mathcal{M}$ is a $K$-continuous chain.
2. $M$ is compatibility prime over $\mathcal{M}$ via $f$.

We often write $M_\delta$ is canonically prime over $\mathcal{M}_\delta = \langle M_i : i < \delta \rangle$ for the formal expression $\text{cpr}(\mathcal{M}_\delta, \mathcal{M}_\delta, f|\delta)$. That is, for brevity we do not mention the specific embeddings $f_{i,j}$ unless they play an active role in the discussion.

**Ch2** For any $K$-continuous chain $\mathcal{M}$ there is a model $M$ and a family of maps $f$ that satisfy $\text{cpr}(\mathcal{M}, M, f)$.

**Ch3** The canonically prime model over an increasing chain $\mathcal{M}$ is unique up to isomorphism over $\mathcal{M}$.

Again the uniqueness axiom is regarded as a desirable property to prove and is not assumed in the general development.

**II.3.2 Definition.** The chain $\langle M_i, f_{i,j} : i < j < \beta \rangle$ is called essentially $K$-continuous if for each limit ordinal $\delta < \beta$, there is a model $M_\delta$ which can be interpolated between the predecessors of $M_\delta$ and $M_\delta$ and is canonically prime over $\mathcal{M}_\delta$.

More formally, we add to the index set a new $\delta'$ for each limit ordinal $\delta$. There exists a system of embeddings $g$ such that if $\alpha$ and $\beta$ are successor ordinals $g_{\alpha,\beta} = f_{\alpha,\beta}$, and for each limit ordinal $\delta$ and each $\alpha < \delta$ $g_{\alpha,\delta}$ factors as $g_{\alpha,\delta'} \circ g_{\delta',\delta}$. Finally $\text{cpr}(\mathcal{M}_\delta, \mathcal{M}_{\delta'}, g|\delta')$.

The next axiom provides a local character for dependence of the prime model over a chain over another model. This is our only basic axiom connecting the independence of an infinite diagram with that of its constituents.

**II.3.3 Axiom L1: Local Dependence for L Diagrams.** If $\langle M_i : i \leq \delta \rangle$ is a $K$-continuous increasing sequence inside $M'$ and for each $i < \delta$, $M_i \downarrow_{M_0} N$ in $M'$ then $M_\delta \downarrow_{M_0} N$ in $M'$. 
CHAPTER II. PRIME MODELS OVER DIAGRAMS

Figure II.3: Local Dependence for L’s
Chapter III

Prime models over small loose trees

In this chapter we extend the existence of prime models over simple diagrams to obtain prime models over more complicated diagrams. Ideally we would show the existence of prime models over an arbitrary independent tree. This project runs into difficulties when considering trees of height greater than $\omega$; as a substitute we show how to obtain prime models over ‘loose trees’ of models indexed by subsets of $\lambda^<\omega$. In a later paper we expect to reduce the discussion of prime models over trees with large height to treatment of loose trees of height $\omega$.

While the original motivation for loose trees was the reduction of problems about tall trees to problems about loose trees of countable height, it turns out that loose trees are closed under several useful operations such as quotient and trees are not. Thus, by passing through loose trees we are able to obtain results about trees that, at least a priori, are otherwise unavailable.

Consider a concrete stable diagram. An initial strategy for building a prime model over $\bar{M}$ is to enumerate $\bar{M}$ say as $M_i$ for $i < \alpha$ and choose a family of models $N_i$ for $i < \alpha$ so that for each $i$, $M_{i+1}$ is independent from $N_i$ over the predecessors of $M_{i+1}$ in the tree and then to take $N_{i+1}$ prime over $M_{i+1}$ and $N_i$. Take canonically prime models over the earlier $N_i$ at limits. The resulting model clearly depends on the order of enumeration. Can one still prove that this model is compatibility prime over the diagram? We show that the answer is yes if the diagram
is a ‘short’ tree (and generalize to allow loose trees). Finite trees are
considered in Section III.1; trees of countable height in Section III.2.
In the case of finite trees we show fairly directly that if a loose tree is
free under one enumeration then it is free under any enumeration. In
the second case we pass to the ostensibly more general notion of locally
free loose tree, show the existence of prime models over such a tree,
and deduce from that the fact that a locally free loose tree is free under
any enumeration.

Assumption: An Adequate Class. We assume in this chapter
axiom groups A and C, axiom D1 and D2 from group D, Ch1, Ch2
from group Ch, and (beginning with III.2.6) L1. We call a class with
these properties an adequate class.

III.1 Free Loose Trees

In this section we show that if $M$ is a finite free loose tree (definitions
follow) of models from an adequate class then there is an absolutely
prime model over $M$. More precisely, we say that $N$ is explicitly prime
over a finite loose tree $M$ if it is the last in a sequence of prime models
over vee’s satisfying certain conditions. We show that $N$ is explicitly
prime over a subdiagram of $M$ then the sequence witnessing this can be
extended to one witnessing the existence of an explicitly prime model
over $M$. This ostensibly technical result is essential for the discussion
of prime models over infinite loose trees in the next section.

III.1.1 Notation. A tree $T$ is a partially ordered set which is isomor-
phic to a subset of $\lambda^{<\omega}$ which is closed under initial segment. We will
often deal directly with this representation. A tree is partially ordered
by containment. If $s, t \in T$ then $s \land t$ denotes the largest common
initial segment of $s$ and $t$ and for any $t$ other than the root, denoted
$\langle \rangle$, $t^-$ denotes the predecessor of $t$.

As in the study of stable diagrams we work only with embeddings
that are inclusions, that is, with concrete diagrams.

III.1.2 Definition. A loose tree of models $M = \{M_t : t \in T\}$ inside
$M$ indexed by a tree $T$ is a collection of $K$-models such that if $t^- = s$
then $M_t \cap M_s \leq M_t$ and $M_t \cap M_s \leq M_s$. 
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Note that any stable diagram indexed by a tree is a loose tree.

We have *not* introduced a loose tree as a kind of partial order. All our index sets are trees in the normal sense. A loose tree of models is loose because of the inclusions amongst the models determined by the indexing.

In defining an independent loose tree (below) we will speak of independence over $M_t \cap M_s$. This use of intersection depends on Axiom C7 (which implies that $M_s \cap M_t \in K$). However, at a cost in complexity of notation the results of this chapter could be obtained through modifying the definition of a loose tree and a free loose tree by replacing $M_t \cap M_s$ by a substructure $M_{t,s} \in K$. This would allow us to extend the definition of loose trees to abstract diagrams. With such an extension the analogy would be closer between the notion of compatibility prime (over a stable embedding of an abstract diagram II.1.10) and the definition III.1.16 of a compatibility prime model over a loose tree.

III.1.3 Definition. An isomorphism of loose trees $f : M \rightarrow M'$ is a family of isomorphisms $f_t$ taking $M_t \in M$ to $M'_t \in M'$ such that for $t \leq s$, $f_t(M_t \cap M_s) \subseteq f_s$. We say the loose tree $M$ can be $K$-embedded in $N$ if there is such an isomorphism between $M$ and a loose tree $M'$ inside $N$.

III.1.4 Definition. A wellordering $\bar{t} = \{t_i : i < \beta\}$, where $|\beta| = |T|$, such that if $t_i$ precedes $t_j$ in $T$ then $i \leq j$ is called an enumeration of the tree $T$.

An enumeration of $T$ induces an enumeration of any $M$ indexed by $T$.

III.1.5 Example. The following observation does not figure in our argument but illustrates one of the subtle differences between a loose tree and a tree of models. The partial order of a tree of models is determined by containments among the models. The partial order of a loose tree is artificially imposed (and thus cannot be ignored). Consider a collection of models $M_0 \subseteq M_1$, $M_0 \subseteq M_2 \subseteq M_3 \subseteq M_4$ with $M_1 \downarrow M_0 M_4$ and $M_1 \cap M_4 = M_0$. As a loose tree of models they can be indexed by the integers less than 5 with the following tree order: 0 is the root, 1 and 2 are incomparable successors of 0, 3 and 4 are incomparable
successors of 2. Clearly \( \langle 0, 1, 2, 3, 4 \rangle \) is an enumeration of the loose tree; so is \( \langle 0, 1, 2, 4, 3 \rangle \).

Again, we could index these models by a tree with 0 as an initial element, 2, 3, 4 as incomparable successors and 1 above 2. Now one can enumerate the tree as \( \langle 0, 2, 3, 4, 1 \rangle \).

With this housekeeping out of the way we can introduce a more important concept.

**III.1.6 Definition.** The loose tree \( M \) indexed by \( T \) is free or independent in \( M \) with respect to an enumeration \( T \) of \( T \) if there exists an ordinal \( \beta \) and a sequence of \( K \)-models \( \langle N_i : i < \beta \rangle \) such that

1. All the \( M_t, N_i \leq M \).
2. \( M_{t_0} = N_0 \).
3. Fix \( k < i \) such that \( t_i = t_k \)
   - If \( 1 \leq i < \omega \) then \( M_t \downarrow_{M_{t_i} \cap M_{t_k}} N_{i-1} \) in \( N_i \).
   - If \( i \geq \omega \) then \( M_t \downarrow_{M_{t_i} \cap M_{t_k}} N_i \) in \( N_{i+1} \).
4. \( \langle N_i : i < \beta \rangle \) is \( K \)-continuous.

Thus \( |\beta| = |T| \) but when \( T \) is infinite the ordinality of \( \beta \) will be \( \text{ord}(T) + 1 \) which may be greater than \( |T| \). If \( M \) is actually a tree the \( M_{t_i} \cap M_{t_k} \) in condition iii) becomes \( M_{t_k} \). We call \( N \) a witnessing sequence for the freeness.

**III.1.7 Remark.** The two conditions in iii) of III.1.6 could be combined if we indexed the \( N_i \) by \( 1 \leq i \leq \beta \). We didn’t make the change to avoid introducing errors in the proofs of the later theorems using finite trees indexed in accordance with the official definition.

A trivial induction shows the following refinement of the definition. We use this observation without comment below.

**III.1.8 Lemma.** If \( M \) is a loose tree in \( M \) indexed by the finite tree \( T \), the witnessing sequence \( \langle N_i : i < \beta \rangle \) can be chosen to satisfy \( N_i \) is prime over \( M_{t_i} \cup N_{i-1} \) inside \( N \).
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With this lemma in mind, suppose the loose tree $M$ indexed by $T$ is free or independent in $M$ with respect to the enumeration $\bar{T}$ of $T$ and there is an initial sequence $\langle M_t : i < j \rangle$ of the enumeration with $M_t \subseteq M_{t+1}$. Then for $i < j$, $N_i$ can be chosen as $M_{t_i}$. This is a likely possibility for a tree, an unlikely one for a properly loose tree of models.

III.1.9 Some generalizations. The following variant may turn out to be necessary. Define a loose tree to be almost free in $N$ with respect to an enumeration $\bar{t}$ if it is free in some extension $N'$ of $N$. Then a compatibility class of $M$ is almost free with respect to $\bar{t}$ if some $M$ is free with respect to $\bar{t}$ in some $N$ in the compatibility class. We will finally show in this section that almost free implies free by showing the explicitly prime model over $M$ inside $N'$ can be chosen inside $N$. This paragraph is analogous to the monotonicity conditions on freeness over a vee.

The following lemma is proved by a straightforward induction on the length of the enumeration using the existence of prime models over independent pairs and canonically prime models over chains.

III.1.10 Lemma. Let $M$ be a loose tree in $M$ enumerated by $\bar{T}$. There is an $N$ and an isomorphic copy of $M$ inside $N$ that is free with respect to the enumeration $\bar{T}$ in $N$.

Note that the isomorphism in Lemma III.1.10 is an isomorphism of loose trees. Example III.1.5 is not free in either enumeration. We want to replace ‘free with respect to an enumeration’ in Lemma III.1.10 by ‘free’; that is to show that if $M$ is free with respect to one enumeration it is free with respect to any enumeration. Then we will show that compatibility prime models exist over free trees. We first handle the case of finite trees. For this we need some combinatorial tools to reduce the main result to a more manageable case.

III.1.11 Definition. Let $\bar{s}$ and $\bar{t}$ be enumerations of $T$. Then, $\bar{s}$ and $\bar{t}$ are close neighbors if for some $i$, $s_i = t_{i+1}$, $t_i = s_{i+1}$, and they agree on all other arguments. $\bar{s}$ and $\bar{t}$ are neighbors if there is a sequence $\bar{t} = \bar{t}^0, \bar{t}^1, \ldots, \bar{t}^k = \bar{s}$ such that for each $j < k$, $\bar{t}^j$ and $\bar{t}^{j+1}$ are close neighbors.
CHAPTER III. PRIME MODELS OVER SMALL LOOSE TREES

The next lemma reduces many problems about the relations between two enumerations to the relation between neighbors and thus by easy induction to relations between close neighbors.

III.1.12 Lemma. If \( \vec{s} \) and \( \vec{t} \) are enumerations of \( T \) then they are neighbors.

Proof. Fix \( \vec{s} \) and \( \vec{t} \). Choose \( k \) maximal such that for some neighbor \( \vec{t}' \) of \( \vec{t} \), \( \vec{s}|k = \vec{t}'|k \). Choose the least \( l \), necessarily greater than \( k \), such that for some such \( \vec{t}' \), \( s_{k+1} = t'_{l} \). Note that \( s_{k+1} = t'_{l} \) is incomparable with \( t'_{l-1} \), \( l - 1 \geq k \). Now if we manufacture \( \vec{t}'' \) from \( \vec{t}' \) by switching \( t'_{l} \) and \( t'_{l-1} \) we contradict the minimality of \( l \).

III.1.13 Remark. Our notion of enumeration corresponds to the concept of ‘linear extension’ in the theory of partial orderings. We are not aware if this result has been proved in that context although it seems likely.

III.1.14 Lemma. If \( T \) is a finite tree and \( M \) is a loose tree indexed by \( T \) which is free inside \( N \) with respect to the enumeration \( \vec{t} \) then \( M \) is free inside \( N \) with respect to any enumeration \( \vec{s} \) of \( M \).

Proof. By Lemma [III.1.12] it suffices to show that if \( M \) is independent with respect to \( \vec{t} \) then \( M \) is independent with respect to the enumeration \( \vec{t} \circ (i, i+1) \). Let \( i = k + 1 \). Since \( \vec{t} \) and \( \vec{s} = \vec{t} \circ (i, i+1) \) are enumerations, \( t_{i} \) and \( t_{i+1} \) are incomparable. We will construct a witnessing sequence \( N'_{i} \) for the independence of \( M \) with respect to \( \vec{s} \) from the given sequence \( N_{i} \) witnessing the independence of \( M \) with respect to \( \vec{t} \). For \( l \leq k \), let \( N'_{l} = N_{l} \). To simplify notation, let \( M^{a} \) denote \( M_{t_{i}} \), \( M^{a} \) denote \( M_{t_{i-1}} \), \( M^{b} \) denote \( M_{t_{i+1}} \), and \( M^{b} \) denote \( M_{t_{i+1}} \). We have \( M^{a} \downarrow M_{a} \cap M_{b} N_{i-1} \) in \( N_{i} \) and \( M^{b} \downarrow M_{b} \cap M_{b} N_{i} \) in \( N_{i+1} \). Since \( M_{b} = M_{t_{i}} \) for some \( l < i \), \( M_{b} \subseteq N_{i-1} \). By monotonicity, \( M^{b} \downarrow M_{b} \cap M_{b} N_{i-1} \) in \( N_{i} \). Now choose \( N'_{i} \leq N_{i+1} \) prime over \( M^{b} \cup N_{i-1} \). By the base extension axiom \( N_{i} \downarrow N_{i-1} N'_{i} \) in \( N_{i+1} \). Since we also have \( M^{a} \downarrow M_{a} \cap M_{b} N_{i-1} \) in \( N_{i} \) and \( M^{a} \subseteq N_{i} \) transitivity of independence, Theorem [II.2.3] yields \( M^{a} \downarrow M_{a} \cap M^{a} N'_{i} \) in \( N_{i+1} \). Since \( N'_{i} \leq N_{i+1} \) and if \( l > i \), \( M_{t_{i}} \downarrow M_{t_{i}} N_{i} \) implies \( M_{t_{i}} \downarrow M_{t_{i}} N'_{i} \) we can let \( N'_{i} = N_{i} \) for \( l > i + 1 \).
III.1. FREE LOOSE TREES

III.1.15 Definition. Henceforth we will say a finite loose tree of models is \textit{free} if it is free under some enumeration.

We want to find an analogous result for certain loose trees indexed by subsets of $\lambda^{<\omega}$. This requires several further concepts. We begin by extending our notion of prime over a stable diagram to free loose trees. Compare the following definition to Definitions II.1.8 and II.1.10.

III.1.16 Definition. The model $M$ is \textit{compatibility prime over the free loose tree} $\underline{M}$ \textit{with respect to the enumeration} $\bar{t}$ via the embedding $f$ of $\underline{M}$ into $M$ if the image of $\underline{M}$ under $f$ is free inside $M$ with respect to $\bar{t}$ and for every $N \in K$ and every $g$ embedding $\underline{M}$ into $N$ with the image of $g$ free inside $N$ such that $M$ and $N$ are compatible over $\underline{M}$ via $f$ and $g$, there is a $\hat{g}$ mapping $M$ into $N$ such that $g = \hat{g} \circ f$.

As in the definition of prime over a stable diagram, we need some ‘freeness’ condition on a loose tree before it makes sense to speak of a ‘prime’ model over it. If we strengthen the definition by allowing $N$ to be arbitrary rather than requiring that $N$ be compatible over $\underline{M}$ with $M$ we call $M$ absolutely prime over $\underline{M}$.

We need a new ‘categorical’ definition for prime over a loose tree because the hidden requirements on composition of maps are much looser when a family is indexed as a loose tree than in our notion of prime model over a stable diagram.

When we restrict to finite loose trees of models we can define absolutely prime models.

III.1.17 Lemma. \textit{If $T$ is a finite tree and} $\underline{M}$ \textit{is a loose tree of models indexed by} $T$ \textit{which is free inside} $N$ \textit{with respect to an enumeration} $\bar{t} = \langle t_i : i < k \rangle$ \textit{then there is an absolutely prime model over the loose tree} $\underline{M}$.

Proof. An easy induction on $|T|$ shows $N_{k-1}$ is the required model.

We christen the result of the last construction.

III.1.18 Definition. i) If $\langle N_i : i < k \rangle$ witnesses that the finite loose tree $\underline{M}$ is free then we say $N_{k-1}$ is \textit{explicitly prime over} $\underline{M}$ (with respect to the embedding map and a specific enumeration).
ii) If \( \langle N_i : i < \beta \rangle \) witnesses that the loose tree \( M \) is free and \( \beta \) is a limit ordinal then we say \( N_\beta \) is explicitly prime over \( M \) (with respect to the embedding map and a specific enumeration) if \( N_\beta \) is canonically prime over \( \langle N_i : i < \beta \rangle \).

Now we invoke the prime models to strengthen the sense in which ‘freeness’ is independent of the enumeration of \( M \). In order to state the result we need several further notations.

**III.1.19 Definition.** i) An *ideal* of a tree \( T \) is a subset \( T_1 \) of \( T \) that is closed under initial segment.

ii) If \( I \) is an ideal of \( T \) then \( T_I \) denotes the (quotient) tree whose elements are \((T - I) \cup \{ \langle \rangle \}\) with the same meets as \( T \) if possible but, if \( x \land y \) in the sense of \( T \) is in \( I \), then \( x \land y = \langle \rangle \) in the sense of \( T_I \).

Any diagram can naturally be restricted to a subset \( X \) of its index set \( T \) by forgetting the models attached elements \( T - X \). We describe several conditions when the restriction \( M|I \) of a (loose) tree \( M \) is a (loose) tree. For loose trees (but not trees) there is a natural complement or quotient structure \( M_I \) to \( M|I \) which we describe in iii) of the next definition.

**III.1.20 Notation.** i) If \( X \subseteq T \) then for any set of models \( M \) indexed by \( T \) there is a natural notion of the restriction \( M|X \) of \( M \) to \( I \): \( M|X = \{ M_t : t \in X \} \).

ii) \( M_I \) denotes some compatibility prime model over \( M|I \).

iii) Suppose \( M = \langle M_t : t \in T \rangle \) is a free loose tree inside \( M \). If \( I \) is an ideal of \( T \) then \( M_I \) denotes \( \langle M_t : t \in T_I \rangle \).

The term \( M_I \) is an abuse of notation since we are choosing one of a number of possible compatibility prime models. If \( I \) is an ideal of \( T \) and \( M \) is a free loose tree of models indexed by \( T \) it is easy to see that the quotient tree is also free. Formally,

**III.1.21 Lemma.** If \( M \) is a free loose tree inside \( M \) and \( I \) is an ideal of \( T \) then \( M_I \) is a free loose tree inside \( M \) indexed by \( T_I \).

Now we consider substitutions for a model in a loose tree.
III.1.22 Notation. Suppose \( M = \langle M_t : t \in T \rangle \) is an indexed family of models. Then \( M(N/s) \) denotes the indexed family of models obtained by replacing \( M_s \) by \( N \). \( M(N_1/s_1, N_2/s_2) \) is the natural extension of this notation to allow two substitutions.

We are interested in a number of substitutions in indexed trees. Most will be approximations to the result of replacing the root in the tree \( M_I \) by the prime model over \( I : M_I(M_I/\langle \rangle) \). The next lemma describes two slightly less simple ways of deriving new loose trees from a given one. The fact that loose trees are obtained by the constructions is an immediate verification; the more important fact that the second construction preserves freeness is proved in Lemma III.1.24.

III.1.23 Lemma. Suppose \( M = \langle M_t : t \in T \rangle \) is a loose tree inside \( M \).

i) Fix \( s \in T \) and a model \( N \leq M \) such that for each \( t \) with \( t^- = s \) or \( s^- = t \), \( N \cap M_t \leq M_t \) and \( N \cap M_t \leq N \). Then \( M' = M(N/s) \) is a loose tree inside \( M \).

ii) Fix \( s, r \in T \) with \( r^- = s \) and \( M_s \leq M_r \). Let \( T' \) denote \( T - \{r\} \). Then \( M' = M|T'(M_r/s) \) is a loose tree inside \( M \).

We have established the notation to state the following lemma.

Lemma III.1.24 (The Omission Lemma). Suppose \( M = \langle M_t : t \in T \rangle \) is finite and free loose tree inside \( M \), \( s, r \in T \) with \( r^- = s \) and \( M_s \leq M_r \). Let \( T' \) denote \( T - \{r\} \). Then \( M' = M|T'(M_r/s) \) is free.

Proof. Let \( \bar{t} \) enumerate \( T \) with \( r = t_{j+1} \) and \( s = t_j \). Suppose \( N \) witnesses the freedom of \( M \) and \( |T| = k \). Now define

\[
\begin{align*}
t'_i & = \begin{cases} 
  t_i & \text{if } i < j \\
  t_{i+1} & \text{if } j \leq i < k - 1
\end{cases} \\
N'_i & = \begin{cases} 
  N_i & \text{if } i < j - 1 \\
  N_{i+1} & \text{if } j - 1 \leq i < k - 1
\end{cases}
\end{align*}
\]

It is straightforward that \( N' \) witnesses the freeness of \( M' \) with respect to the enumeration \( \bar{t}' \).
CHAPTER III. PRIME MODELS OVER SMALL LOOSE TREES

The following theorem asserts that the freeness of a finite loose tree is independent not only of the enumeration but for a given enumeration of the choice of the prime models $N_i$ witnessing the freeness. This theorem allows us to decompose $M$ into $M|I$ and $M_I$. An advantage of loose trees in performing this construction is that they allow the mixing of the models in the tree and the witnessing sequence.

III.1.25 Theorem. Suppose $M$ is a free loose tree of models inside $M$ indexed by the finite tree $T$ and $I$ is an ideal of $T$. Suppose also that $N$ is explicitly prime over $M|I$ for an enumeration $t$ of $I$. Then $M_I(N/\langle \rangle)$ is free inside $M$.

Proof. Suppose $|I| = j$. Extend the given enumeration of $I$ to an enumeration of $T$ and also denote the extension by $T$. Let $N$ witness that $M|I$ is free. Noting that any initial segment of an enumeration is an ideal, let $T_i$ be the ideal composed of the first $i$ elements in the enumeration. Let $M'_i = M_{T_i}(N_i/\langle \rangle)$. We finish by showing by induction on $i < |I|$ that $M'_i$ is free inside $M$.

If $i = 0$, $M'_i = M$ and the result is clear. Suppose we know $M'_i$ is free and consider $M'_{i+1}$. In $T_{t_i}$, $t_{i+1}$ is an immediate successor of $\langle \rangle$. Lemma III.1.26 below shows $M'_i(N_{i+1}/t_{i+1})$ is free inside $M$. This implies $M'_{i+1}$ is free by the omission lemma and the induction hypothesis. Thus, it remains only to prove Lemma III.1.26.

Note that any finite tree can be enumerated so that $M_{t_i+1} \leq M_i$ for any $i$.

III.1.26 Lemma. Suppose $M$ is a free loose tree of models inside $M$ indexed by the finite tree $T$. Let $N \leq M$ and suppose $N$ is prime over $M_v \cup M_v^-$ for some $v \in T$. Then $M'(N/v)$ is free inside $M$.

Proof. Fix an enumeration $\bar{t}$ of $M$ with $t_j = v$. Suppose that $N = \langle N_i : i < |T| \rangle$ witnesses the freedom of $M$. Before constructing the sequence $N'_i$ which will witness the freedom of the new tree we need an auxiliary sequence. We define for $j \leq i < |T|$ an increasing chain of models $N_i^*$ beginning with $N_j^* = M_v = M_{t_j}$ such that

i) each $N_i^* \leq N_i$
### III.1. FREE LOOSE TREES

| $N_j$ | $N_{j+1}$ |
|-------|-----------|
| $M_v$ | $N_{j+1}^*$ |
| $M_{t_{j+1}} \cap M_v$ | $M_{t_{j+1}}$ |

Figure III.1: The base step

- ii) $M_{t_{i+1}} \downarrow_{M_{t_{i+1}} \cap M_{t_{i+1}}} N_i^*$ inside $N_{i+1}^*$.
- iii) $N_{i+1}^* \downarrow_{N_i^*} N_i$ inside $N_{i+1}$,
- iv) $N_{i+1}^*$ is prime over $M_{t_{i+1}} \cup N_i^*$.

**Base Step:** First note setting $N_j^* = M_j$ satisfies conditions ii) and iii). Choosing $N_{j+1}^*$ to satisfy iv), we must verify iii). This follows by the base extension axiom applied as in Figure III.1. Recall $M_j = M_v = N_j^*$.

**Induction Step:** Suppose that for some $k \geq j$ we have chosen $N_k^*$ to satisfy ii) and iii) (with $i = k - 1$). Choose $N_{k+1}^*$ inside $N_{k+1}$ prime over $M_{t_{k+1}} \cup N_k^*$ to satisfy i) and iv). Now $N_{k+1}^*$ satisfies ii) by monotonicity and iii):

$$N_{k+1}^* \downarrow_{N_k^*} N_k \text{ in } N_{k+1}$$

follows from the base extension axiom as in the induction step diagram, Figure III.2.

This completes the construction of the $N_i^*$.

Let $u$ denote $v^-$; note that since $N_j^* = M_v$ we have $N_{j-1} \downarrow_{M_v \cap M_u} N_j^*$. With this as the base an easy induction on $l$ for $j \leq l \leq k$ shows $N_{l-1} \downarrow_{M_v \cap M_u} N_l^*$. We will rely on the case $l = k$.

We have from the free enumeration that $M_v \downarrow_{M_v \cap M_u} N_{j-1}$. The base extension axiom then gives, since $N$ is prime over $M_v \cup M_v^-$, $N \downarrow_{M_v} N_{j-1}$. 
Choose $M'_v \subseteq N_k$ prime over $\mathcal{N}\cup\mathcal{N}_{j-1}$. By transitivity of primeness, Lemma II.2.4, $M'_v$ is prime over $M_v \cup \mathcal{N}_{j-1}$.

From the base extension axiom and the conclusions of the last two paragraphs, we deduce $M'_v \downarrow M, N^*_k$. This allows us to perform the $i = j$ step of the following construction. Choose $N'_i$ so that

$$N'_i = \begin{cases} N_i & \text{if } i < j \\ M'_v & \text{if } i = j \end{cases}$$

and for $j < i$ so that $N'_i \leq M$ is prime over $N^*_i \cup N'_{i-1}$ and $N'_i \downarrow N^*_i \cup N'_k$. The choice of the $N'_i$ for $i > j$ is a straightforward induction.

Now to complete the proof we must observe that the $N'_i$ witness the freeness of $M'_v$. That is we must show that $N'_{i+1}$ is prime over $M_{i+1} \cup N'_i$. This follows from transitivity of primeness, Lemma II.2.4, since for each $i$ we have $N^*_i \cup N'_{i+1}$ is prime over $M_{i+1} \cup N^*_i$ and $N'_{i+1}$ is prime over $N^*_i \cup N'_i$.

We restate Theorem III.1.25 in a more applicable form.

### III.1.27 Corollary.

Suppose $\mathcal{M}$ is a free loose tree of models inside $\mathcal{M}$ indexed by the finite tree $\mathcal{T}$ and $\mathcal{I}$ is an ideal of $\mathcal{T}$. If $\bar{T}$ is an enumeration of $\mathcal{I}$ and $\mathcal{N}$ witnesses that $\mathcal{M}|\mathcal{I}$ is free inside $\mathcal{M}$ then $\mathcal{N}$ can be extended to a sequence witnessing that $\mathcal{M}$ is free inside $\mathcal{M}$.

We have shown that if $\mathcal{K}$ is an adequate class and $\mathcal{M}$ is a finite (loose) tree $\mathcal{M}$ which is free inside $\mathcal{N}$ under some enumeration then $\mathcal{M}$
is free under any enumeration and there is an absolutely prime model over $M$. With more difficulty we showed that if $I$ is an ideal of $M$ and $N_0$ is explicitly prime over $I$ then there is an sequence defining an explicitly prime model over $M$ which includes $N_0$.

### III.2 Locally Free Loose Trees

We now extend our analysis to infinite trees. Recall that in the absence of the monster model we can only speak of a diagram being free when we have an embedding into an ambient model in mind. And only then can we discuss the possibility of any sort of ‘prime’ model over the diagram. This freeness can be verified by an enumeration of the loose tree and if the index tree is finite the freeness is independent of the enumeration. We do not at first claim so much for an infinite tree. Rather we introduce a notion of a locally free (loose) tree that has the following properties. If $M$ is free under some enumeration then $M$ is locally free. The property of being locally free does not depend on the enumeration. We will establish the existence of compatibility prime models over locally free loose trees (which have height at most $\omega$). From this we deduce that if such a tree is free is under one enumeration then it is free under any enumeration.

#### III.2.1 Definition

A loose tree of models $M$ is **locally free in** $N$ if $M = \langle M_t : t \in T \rangle$ is contained in $N$ and for every finite subtree $T_1 \subseteq T$ the finite loose tree $M|T_1 = \langle M_t : t \in T_1 \rangle$ is free inside $N$.

#### III.2.2 Remark

This definition relies on our restriction to subtrees of $\lambda^{<\omega}$. Since subtrees are closed under predecessor trees of greater height cannot be covered by finite subtrees. Thus, when we deal with trees of greater height we will modify the definition of locally free but leave the meaning the same on the low trees considered here.

The next proposition is obvious.

#### III.2.3 Proposition

If the loose tree $M$ is free inside for $N$ for some enumeration then $M$ is locally free inside $N$.

Now we define ‘prime’ models over locally free loose trees. Definition III.1.10 is almost a special case of this. (Ostensibly, a map could take a
free tree to a locally free tree so prime over locally free is more restrictive than prime over free.)

**III.2.4 Definition.** The model $M$ is *compatibility prime over the locally free loose tree $M$* via the embedding $f$ of $M$ into $M$ if the image of $M$ under $f$ is locally free inside $M$ and for every $N \in K$ and $g$ embedding $M$ into $N$ with the image of $g$ locally free inside $N$ such that $M$ and $N$ are compatible over $M$ via $f$ and $g$ there is a $\hat{g}$ mapping $M$ into $N$ such that $g = \hat{g} \circ f$.

The name of this notion is unconscionably long so we resort to the following abbreviation: $\text{LFP}(M, M, f)$. We omit the $f$ if $M$ is concretely realized in $M$ (whence $f$ is family of inclusions.)

We are going to show that if a loose tree $M$ is locally free there is a ‘prime’ model over the loose tree $M$. We need the following notation (extending III.1.19 and III.1.22) for trees obtained from a representation of given tree as a union of (smaller) trees to carefully state and prove this proposition.

**III.2.5 Notation.** Suppose $T = \bigcup_{\alpha < \lambda} I_\alpha$ where the $I_\alpha$ form an increasing continuous sequence of ideals in $T$.

i) $M_\alpha$, denotes the quotient of $M$ indexed by the quotient tree $\{ t : t \in T_{I_\alpha} \}$.

ii) For $\beta < \alpha$, $M_{\alpha, \beta}$ denotes the quotient of $M|_{(I_\alpha)}$ by $I_\beta$. That is, $M_{\alpha, \beta}$ is indexed by the quotient tree $\{ t : t \in (I_\alpha)_{I_\beta} \}$ (Remember $I_\alpha$ is a subtree of $T$ so this is a natural extension of our previous notation III.1.22).

Thus $M_\alpha$ just abbreviates $M_{I_\alpha}$. Although $M_\alpha$ is a quotient of $M$ and $M_{\alpha, \beta}$ is a quotient of $M|_{(I_\alpha)}$, as sets of models $M_\alpha$ and $M_{\alpha, \beta}$ are contained in $M$.

Here is the main result of this section. Condition 1 is the result we really want (see Conclusion III.2.10) but to establish it we must prove Conditions 1 and 2 by a simultaneous induction. A first try to prove this theorem by induction would decompose a tree $T$ as a union of properly smaller ideals and then choose a ‘prime’ model over each of these ideals and take the canonically prime model over this chain...
III.  LOCALLY FREE LOOSE TREES

of models. The difficulty is to guarantee that the sequence of models forms a chain. The double induction accomplishes this end.

III.2.6 Theorem. Suppose \( \mathcal{M} \) is a locally free loose tree indexed by \( T_0 \) in \( N \). Conditions 1) and 2) will be proved by simultaneous induction on \( \lambda \).

Condition 1. If \( |T_0| = \lambda \) then there is a compatibility prime model for locally free loose trees over \( \mathcal{M} \). That is, there is an \( N \) satisfying \( \text{LFP}(\mathcal{M}, N) \).

Condition 2. If \( I \) is an ideal of \( T_0 \) and \( |I| = \lambda \) then \( \mathcal{M}_I(\mathcal{M}_I/\langle \rangle) \) is locally free inside \( \mathcal{M} \).

Note that the cardinality of \( T_0 \) is not bounded in condition 2. The following lemma provides most of the technicalities of the proof.

Lemma III.2.7 (The key lemma). Suppose \( \mathcal{M} \) is a free loose tree in \( N' \) indexed by the finite tree \( T \) and suppose \( s \) is a minimal (\( \neq \langle \rangle \)) element of \( T \). Suppose also \( \langle M^\alpha_s : \alpha \leq \delta \rangle \) is an increasing \( K \)-continuous sequence satisfying the following conditions.

i) For each \( \alpha < \delta \), \( M^\alpha = M(M^\alpha_s/s) \) is a free tree inside \( N' \).

ii) For each \( \alpha < \delta \), \( M_{\langle \rangle} \leq M^\alpha_s \).

iii) If \( t = s \) then \( M_t \downarrow_{M_t \cap M^\alpha_s} M^\alpha_s \).

Then \( \mathcal{M}(M^\delta_s/s) \) is free inside \( N' \).

Proof. Fix an enumeration \( T \) of \( T \) such that the elements not in the cone with vertex \( s \) come first. Suppose \( s = t_k \) and \( |T| = n \). Let \( N^\alpha \) witness that \( M^\alpha \) is free. By Corollary III.1.27 we may assume that \( N^\alpha_{t_i} = N^0_k \) if \( i < k \). (This use of Corollary III.1.27 only simplifies notation; the later use is essential.) In order to discuss uniformly the trees \( \hat{M}^\alpha \), we refer below to models \( M^\alpha_t \). Unless \( t = s \), \( M^\alpha_t = M^0_t = M_t \).

 Expand the tree \( T \) to \( \hat{T} \) by adding a new element \( r \) with \( r^- = s \) but with no elements above \( r \). We define for each \( \alpha \leq \delta \) a loose tree of models \( \hat{M}^\alpha \) and a sequence of models \( \hat{N}^\alpha \).

\[
\hat{M}^\alpha_x = \begin{cases} 
M^0_x & \text{if } x \in T \\
M^\alpha_s & \text{if } x = r
\end{cases}
\]
\[ \hat{N}_i^\alpha = \begin{cases} N_i^0 & \text{if } i \leq k \\ N_{i-1}^\alpha & \text{if } i > k \end{cases} \]

Consider the following enumeration \( u \) of \( \hat{T} \) which we obtain from \( t \).
\[
u_i = \begin{cases} t_i & \text{if } i \leq k \\ r & \text{if } i = k + 1 \\ t_{i-1} & \text{if } k + 1 < i \leq n - 1 \end{cases}
\]

Recall that if \( i \neq k \), \( M_i^0 = M_i^\alpha \). Note that both \( M_i^0 \) and \( M_s^\alpha \) occur in \( \hat{M}^\alpha \) (\( M_s^\alpha = M_s^0 \), \( \hat{M}_r^\alpha = \hat{M}_{u_{k+1}}^\alpha = M_s^\alpha \)). We will show below that for each \( \alpha \), \( \hat{N}^\alpha \) witnesses that \( \hat{M}^\alpha \) is free inside \( N' \). Assuming this fact we now complete the proof of the lemma.

Since we proved that the freeness of a loose tree does not depend on the enumeration, each \( \hat{M}^\alpha \) is also free by an enumeration that places \( r \) last. That is \( \hat{M}^\alpha \) is free with respect to the enumeration \( \nu \) defined as follows.
\[
u_i = \begin{cases} t_i & \text{if } i \leq n - 1 \\ r & \text{if } i = n \end{cases}
\]

Let \( I = \{v_i : i < n\} \). By Corollary III.1.27 any sequence, in particular \( \hat{N}^0 = \langle N_0^0 \ldots N_{n-1}^0 \rangle \), which witnesses the freeness of \( \hat{M}^\alpha \mid I \) can be extended to a sequence witnessing the freeness of \( \hat{M}^\alpha \). Let \( N \) denote \( N_{n-1}^0 \). (It was to make the choice of \( N \) independent of \( \alpha \) that we needed to prove Theorem III.1.23 and Corollary III.1.27.) Now for each \( \alpha \) since \( \hat{M}^\alpha \) is free with respect to \( \nu \) we have
\[
N \downarrow_{\hat{M}_r^\alpha \cap \hat{M}_{u_{k+1}}^\alpha} \hat{M}_r^\alpha
\]
inside \( \hat{N}_n^\alpha \). That is,
\[
N \downarrow_{M_s^\alpha \cap M_s^0} M_s^\alpha
\]
inside some \( N_n^\alpha \). This implies and since \( M_s^0 \subseteq M_s^\alpha \) that for each \( \alpha \)
\[
N \downarrow_{M_s^0} M_s^\alpha
\]
inside \( N' \). Therefore, by L1,
\[
N \downarrow_{M_s^0} M_s^\delta
\]
inside $N'$. Choosing $N$, prime over $N \cup M_s^\delta$, we finish the proof of the key lemma from the assumption.

We are left with verifying the assumption; that is, showing that for each $\alpha$, $\hat{N}_i^\alpha$ witnesses that $M_s^\alpha$ is free inside $N'$ with respect to $\pi$. We will rely on the following fact. For each $\alpha$ and each $i < n$

$$M_{t_i}^\alpha \downarrow_{M_i^\alpha \cap M_{t_i - 1}^\alpha} N_{t_i - 1}^\alpha \text{ inside } N_i^\alpha. \quad (III.1)$$

We must show that for each $i < n + 1$,

$$M_{s_i}^\alpha \downarrow_{M_i^\alpha \cap M_{s_i - 1}^\alpha} \hat{N}_{t_i - 1}^\alpha \text{ inside } \hat{N}_i^\alpha. \quad (III.2)$$

For $i \leq k$ this is an immediate translation of property $[III.1]$ (with $\alpha = 0$). For $i = k + 1$ property $[III.2]$ translates to

$$M_s^\alpha \downarrow_{M_s^\alpha \cap M_s^\alpha} N_k^0 \text{ inside } N_k^0. \quad (III.3)$$

since $\hat{M}_s^\alpha = M_s^\alpha = M_0^s = M_s^0$. Remember that $N_k^0$ is prime over $N_{k-1}^0 \cup M_s^0$ and since $N_k^0 = N_k^\alpha$, $N_k^0$ is prime over $N_{k-1}^0 \cup M_s^\alpha$. Moreover, $M_s^\alpha \cap M_s^\alpha = M_s^\alpha \cap M_s^\alpha = M_s^\alpha \cap M_s^\alpha = M_s^\alpha \cap M_s^\alpha$. Now property $[III.3]$ follows by base extension from

$$M_s^\alpha \downarrow_{M_s^\alpha \cap M_s^\alpha} N_k^0 \text{ inside } N_k^0. \quad (III.4)$$

which is an instance of property $[III.1]$.

For $i = k + 2$, since $u_{k+2} = t_{k+1} = t_k = s$ property $[III.2]$ becomes

$$M_{t_{k+1}}^\alpha \downarrow_{M_{t_{k+1}}^\alpha \cap M_{s}^\alpha} N_k^\alpha \text{ inside } N_{k+1}^\alpha. \quad (III.5)$$

Now by assumption iii) of the Lemma, $M_{t_{k+1}}^\alpha \downarrow_{M_{t_{k+1}}^\alpha \cap M_{s}^\alpha} M_s^\alpha$ so we can deduce property $[III.3]$ from property $[III.1]$ and transitivity of independence (Theorem $[II.2.3]$).

For $i > k + 2$, property $[III.2]$ translates to

$$M_{t_i - 1}^\alpha \downarrow_{M_{t_i - 1}^\alpha \cap M_{t_i - 2}^\alpha} N_{t_i - 2}^\alpha \text{ inside } N_i^\alpha. \quad (III.6)$$
which follows immediately from property III.1.

We will apply the following consequence of this lemma (with $I_\delta$ a proper ideal of $T$) in the main proof.

**III.2.8 Corollary.** Suppose $M$ is a locally free loose tree in some $N$ which is indexed by $T$ and $\langle I_\alpha : \alpha \leq \delta \rangle$ is a continuous increasing sequence of ideals in $T$. Suppose further that there is a $K$-continuous sequence $\langle N_\alpha : \alpha \leq \delta \rangle$ such that for each $\alpha < \delta$, $\bigcup_{t \in I_\alpha} M_t \subseteq N_\alpha \leq N$ and $M_\alpha(N_\alpha/\langle \rangle)$ is a locally free loose tree inside $N$. Then $M_\delta(N_\delta/\langle \rangle)$ is a locally free loose tree inside $N$.

**Proof.** Without loss of generality, (since we are trying to establish local freeness), $T - I_\delta$ is finite. Let $\{t_i : i < k\}$ be the minimal elements of $T - I_\delta$ and let $s_i \in I_\delta$ be the predecessor of $t_i$ in $I_\delta$. Again, without loss of generality, we may assume all the $s_i$ are in $I_0$. For each $\alpha \leq \delta$, we define a loose tree $\tilde{M}_\alpha$ (not following our previous conventions) as follows. $\tilde{M}_\alpha$ is indexed by $(T - I_\delta) \cup \{s, \langle \rangle\}$ where $s$ is interpolated above $\langle \rangle$ and below each member of $T - I_\delta$. $\tilde{M}_\alpha^t = M_t$ if $t \in (T - I_\delta)$, $N_0$ if $t = \langle \rangle$ and $N_\alpha$ if $t = s$. $\tilde{M}_\alpha^t$ is a free loose tree inside $N$. (Just enumerate $\langle \rangle$, $s$ and then the rest using $M_\alpha(N_\alpha/\langle \rangle)$ is a locally free loose tree inside $N$.) By the choice of the $t_i$ and $s_i$, and the local freeness of $M_\alpha(N_\alpha/\langle \rangle)$ inside $N$ we verify the third assumption of Lemma III.2.7. In the present context (since $M_s \subseteq N_0$ and $M_t \cap M_s \leq N_0$) it translates as for each $i < k$, $M_t \downarrow_{M_t \cap N_0} N_\alpha$. Now by Lemma III.2.7 $\tilde{M}_\delta^t$ is free. Applying the omission lemma to $\tilde{M}_\delta$ we conclude that $M_\delta(N_\delta/\langle \rangle)$ is a locally free loose tree inside $N$.

Now we continue with the proof of the main theorem (Theorem III.2.6).

**III.2.9 Proof of III.2.6.** We assume by induction that Condition 1 of the theorem holds for any tree with cardinality less than $\lambda$ and Condition 2 of the theorem holds for any tree and any ideal of cardinality less than $\lambda$.

Fix a subtree $T_1 \subseteq T_0$ with $|T_1| = \lambda$. We will describe a construction relative to $T_1$. Then to verify each of conditions 1 and 2 we use a different choice of $T_1$. 
As in [III.2.3], suppose $T_1 = \bigcup_{\alpha < \lambda} I_\alpha$ where the $I_\alpha$ are an increasing continuous sequence of ideals in $T_1$ and $|I_\alpha| < \lambda$. We define by induction on $\alpha < \lambda$, an increasing $K$-continuous sequence $\langle N_\alpha : \alpha < \lambda \rangle$ of $K$-submodels of $N$ satisfying the following conditions.

i) If $\alpha = \beta + 1$ then $N_\alpha$ is prime for loose trees over $M_{\alpha,\beta}(N_\beta/\langle \rangle)$.

ii) $M_\alpha(N_\alpha/\langle \rangle)$ is locally free.

Recall that $M_{\alpha,\beta}$ denotes the quotient of $M(I_\alpha)$ by $I_\beta$. Note that $M_\alpha$ is indexed by $(T_0 - I_\alpha) \cup \{\langle \rangle\}$ (not $(T_1 - I_\alpha) \cup \{\langle \rangle\}$). This is crucial for the verification of Condition 2.

There are three cases in the construction.

$\alpha = 0$. Condition i) doesn’t apply. Condition ii) holds by the induction hypothesis applied to Condition 2) of the main theorem.

$\alpha$ is a limit ordinal. Let $N_\alpha$ be canonically prime over $\langle N_\beta : \beta < \alpha \rangle$. Condition i) does not apply and condition ii) is immediate from the last result: Corollary [III.2.8].

$\alpha = \gamma + 1$. Applying condition ii) to $\gamma$, $M_\alpha(N_\alpha/\langle \rangle)$ is locally free in $N$. So the subdiagram $M_{\alpha,\gamma}$ is locally free in $N$. So by induction applied to Condition 1) of the main theorem with $I_\alpha$ as $I$ there is an $N_\alpha$ satisfying condition i). That is, $N_\alpha$ is prime for loose trees over $M_{\alpha,\gamma}(N_\gamma/\langle \rangle)$. In particular, note $N_\gamma$ is embedded into $N_\alpha$. Applying the induction hypothesis for Condition 2 (with $T_0 = M_{I_\alpha}$) $M_\alpha(N_\alpha/\langle \rangle)$ is locally free and we satisfy condition ii).

This completes the construction.

To see that Condition 1 holds of $T_0$, take $T_1$ as $T_0$. We must show that if $N_\lambda$ is taken canonically prime over $\langle N_\beta : \beta < \lambda \rangle$ then $N_\lambda$ satisfies LFP($N, N_\lambda$). Suppose $f = \langle f_t : t \in T \rangle$ maps the loose tree $M$ isomorphically to a loose tree $M'$ which is locally free in some $M'$. We must extend $f$ to an embedding of $N_\lambda$ into $M'$.

For this, repeat the preceding argument constructing a sequence $N_\alpha'$ for $\alpha < \lambda$ inside $M'$ which satisfy conditions i) and ii) (for $M'$) and simultaneously construct a continuous increasing sequence of maps $g_\alpha$
for \( \alpha < \delta \) which map \( N_\alpha \) isomorphically onto \( N'_\alpha \) and such that \( g_\alpha \) extends each \( f_t \) with \( t \in I_\alpha \). We conclude Condition 1.

For Condition 2 we must show that if \( I \) is an ideal of \( T \) and \( |I| = \lambda \) then \( M_I \) is locally free. Applying the construction above with \( I \) as \( T_1 \), we have constructed the \( N_\alpha \) so that \( M_\alpha(N_\alpha/\{\}) \) is locally free. Thus, by Corollary \[ \text{III.2.8} \] \( M_\alpha(N_\delta/\{\}) \) is locally free and we can conclude Condition 2.

III.2.10 Conclusion. Let \( K \) be an adequate class and suppose \( M \) is a loose tree of \( K \)-models indexed by a subtree of \( \lambda^{<\omega} \). If \( M \) is locally free then there is a compatibility prime model \( M \) over \( M \) (i.e. \( \text{LFP}(M, M, f) \), see \[ \text{III.2.4} \])

In fact, if \( M \) (indexed by a subtree of \( \lambda^{<\omega} \)) is free under one enumeration then it is free under any enumeration. For, by Proposition \[ \text{III.2.3} \] it is locally free. Let \( \overline{f} = \langle t_i; i < |T| \rangle \) be an arbitrary enumeration of \( M \). For each \( \gamma < |T| \) let \( I_\gamma \) be the ideal containing each \( t_i \) with \( i < \gamma \). Now construct a family of models \( \langle N_\gamma : \gamma < |T| \rangle \) as in the proof of Theorem \[ \text{III.2.6} \]. By the basic properties of independence (using \( L_1 \) at limit stages) these models witness that \( M \) is free under the given enumeration.

We have shown that if \( M \) is indexed by a subtree of \( \lambda^{<\omega} \) and if \( M \) is free with respect to some enumeration then there is a compatibility prime model over \( M \). In \[ 3 \] we will show that if \( K \) is at all manageable then there is only one compatibility class over \( M \). We are attempting to show that models in manageable classes are 'tree-decomposable' by trees of small height. The existence of such prime models is an essential step in this program.
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