IF ARCHIMEDES WOULD HAVE KNOWN FUNCTIONS ...

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ABSTRACT. Could calculus on graphs have emerged by the time of Archimedes, if function, graph theory and matrix concepts were available 2300 years ago?

1. SINGLE VARIABLE CALCULUS

Calculus on integers deals with functions $f(x)$ like $f(x) = x^2$. The difference $Df(x) = f'(x) = f(x+1) - f(x) = 2x + 1$ as well as the sum $Sf(x) = \sum_{k=0}^{x-1} f(k)$ with the understanding $Sf(0) = 0, Sf(-1) = f(-1)$ are functions again. We call $Df$ the derivative and $Sf$ the integral. The identities $DSf(x) = f(x)$ and $SDf(x) = f(x) - f(0)$ are the fundamental theorem of calculus. Linking sums and differences allows to compute sums (which is difficult in general) by studying differences (which is easy in general). Studying derivatives of basic functions like $x^n, \exp(a \cdot x)$ will allow to sum such functions. As operators, $Xf(x) = xsf(x)$ and $Df(x) = [s, f]$ where $sf(x) = f(x+1)$ and $s^*f(x) = f(x-1)$ are translations. We have $1x = x, Xx = x(x-1), X^2x = x(x-1)(x-2)$. The derivative operator $Df(x) = (f(x+1) - f(x))s$ satisfies the Leibniz product rule $D(fg) = (Df)g + (Dg)f = Dfg + f^*Dg$. Since momentum $P = iD$ satisfies the anti-commutation relation $[X, P] = i$ we have quantum calculus. The polynomials $[x]^n = X^n x$ satisfy $D[x]^n = n[x]^{n-1}$ and $X^n X^m = X^{n+m}$. The exponential $\exp(a \cdot x) = (1 + a)^x$ satisfies $D \exp(a \cdot x) = a \exp(a \cdot x)$ and $\exp(a \cdot (x + y)) = \exp(a \cdot x) \exp(a \cdot y)$. Define $\sin(a \cdot x)$ and $\cos(a \cdot x)$ as the real and imaginary part of $\exp(ia \cdot x)$. From $De^{ax} = ae^{ax}$, we deduce $D \sin(a \cdot x) = a \cos(a \cdot x)$ and $D \cos(a \cdot x) = -a \sin(a \cdot x)$. Since exp is monotone, its inverse log is defined. Define the reciprocal $X^{-1}$ by $X^{-1} X^1 = D \log(x)$ and $X^{-k} = (X^{-1})^k$. We check identities like $\log(x^y) = \log(x) + \log(y)$ from which $\log(1) = 0$ follows. The Taylor formula $f(x) = \sum_{k=0}^{\infty} D^k f(0) X^k/k!$ is due to Newton and Gregory and can be shown by induction or by noting that $f(x + t)$ satisfies the transport equation $D_t f = Df$ which is solved by $f(t) = \exp(Dt)f(0)$. The Taylor expansion gives $\exp(x) = \sum_{k=0}^{\infty} X^k/k!$. We also get from $\frac{1}{1-x} = \sum_{k=0}^{\infty} X^k$ that $\log(1-x) = -\sum_{k=1}^{\infty} X^k/k$. The functions $c \sin(a \cdot x) + d \cos(a \cdot x)$ solve the harmonic oscillator equation $D^2f = -a^2f$. The equation $Df = g$ is solved by the anti-derivative $f = Sg(x) + c$. The wave equation $D^2_t f = D^2 f$ can be written as $(D_t - D)(D_t + D)f = 0$ which thus has the solution $f(t)$ given as a linear combination of $\exp(Dt)f(x) = f(x + t)$ and $\exp(-D_t)f(x) = f(x - t)$. We replace now integers with a graph, the derivative $iD$ with a matrix $D$, functions by forms, anti-derivatives by $D^{-1}f$, and Taylor with Feynman.
2. Multivariable calculus

Space is a finite simple graph \( G = (V,E) \). A subgraph \( H \) of \( G \) is a curve if every unit sphere of a vertex \( v \) in \( H \) consists of 1 or 2 disconnected points. A subgraph \( H \) is a surface if every unit sphere is either an interval graph or a circular graph \( C_n \) with \( n \geq 4 \). Let \( \mathcal{G}_k \) denote the set of \( K_{k+1} \) subgraphs of \( G \) so that \( G_0 = V \) and \( G_1 = E \) and \( \mathcal{G}_2 \) is the set of triangles in \( G \). Intersection relations render each \( \mathcal{G}_k \) a graph too. The linear space \( \Omega_k \) of anti-symmetric functions on \( \mathcal{G}_k \) is the set of \( k \)-forms. For \( k = 0 \), we have scalar functions. To fix a basis, equip simplices like edges and triangles with orientations. Let \( \partial \) denote the exterior derivative \( \Omega_k \to \Omega_{k+1} \).

If \( \delta H \) be the boundary of \( H \), then \( df(x) = f(\delta x) \) for any \( x \in \mathcal{G}_k \). The boundary of a curve \( C \) is empty or consists of two points. The boundary of a surface \( S \) is a finite set of closed curves. A closed surface \( S \) is a discretisation of a 2-dimensional classical surface as we can build an atlas from wheel subgraphs. An orientation of \( S \) a choice of orientations on triangles in \( S \) such that the boundary orientation of two intersecting triangles match. If \( F \) is a 2-form and \( S \) is an orientable surface, define \( \int_S F \, dS = \sum_{x \in \mathcal{G}_2(S)} F(x) \). An orientation of a curve \( C \) is a choice of orientations on edges in \( C \) such that the boundary orientation of two adjacent edges correspond. If \( f \) is a 1-form and \( C \) is orientable, denote by \( \int_C F \, ds \) the sum \( \sum_{x \in \mathcal{G}_1(C)} F(x) \).

For a 0-form, write \( \text{grad}(f) = df \) and for a 1-form \( F \), write \( \text{curl}(F) = dF \). If \( S \) is a surface and \( F \) is a 1-form, the theorem of Stokes tells that \( \int_S dF = \int_{\partial S} F \). It is proved by induction using that for a single simplex, it is the definition. If \( H \) is a 1-dimensional graph, and \( F \) is a function, then \( \int_H dF = \int_{\delta H} F \) is the fundamental theorem of line integrals. A graph is a solid, if every unit sphere is either a triangulation of the 2-sphere or a triangulation of a disc. Points for which the sphere is a disc form the boundary. If \( F \) is a 2-form and \( H \) is a solid, then \( \int_H dF = \int_{\delta H} F \) is Gauss theorem. The divergence is the adjoint \( d^*: \Omega_1 \to \Omega_0 \). In school calculus, 1-forms and 2-forms are often identified as “vector fields” and 3-forms on the solid \( H \) is treated as scalar functions. Gauss is now the divergence theorem. Define the Dirac matrix \( D = d + d^* \) and the form Laplacian \( L = dd^* + d^*d \). Restricted to scalar functions, it is \( d^*df = \text{div} (\text{grad} f) \) which is equal to \( A-B \), where \( A \) is the adjacency matrix and \( B \) the diagonal degree matrix of the graph. The Schrödinger equation \( f' = iDf \) is solved by \( f(t) = e^{itD}f(0) = f(0) + tf(0) + D^2t^2f(0)/2! + \ldots \). It is a multi-dimensional Taylor equation, because by matrix multiplication, the term \( (t^nD^n)f \) is the sum over all possible paths \( \gamma \) of length \( n \) where each step \( k \) is multiplied with \( tf(\gamma(k)) \). The solution of the Schrödinger equation \( f(t,y) \) is now the average over all possible paths from \( x \) to \( y \). This Feynman-Kac formula holds for general \( k \)-forms. Given an initial form \( f \), it describes \( f(t) \). The heat equation \( f' = -Lf \) has the solution \( e^{-Lt}f(0) \), the heat flow. Because \( L \) leaves \( k \)-forms invariant, the Feynman-Kac formula deals with paths which stay on the graph \( \mathcal{G}_k \). As any linear differential equation, it can be solved using eigenfunctions of \( L \). If \( G = C_n \), it is a discrete Fourier basis. The wave equation \( f'' = -Lf \) can be written as \( (\partial_t - iD)(\partial_t + iD)f = 0 \) which shows that \( f = \cos(Dt)f(0) + \sin(Dt)D^{-1}f'(0) = e^{iDt} \psi \) with \( \psi = f(0) - iD^{-1}f'(0) \). Feynman-Kac for a \( k \)-form expresses \( \psi(t,x) \) as the expectation \( \mathbb{E}[\sum_{x \in \gamma} t^n(\gamma)f(0,x)] \) on a probability space of finite paths \( \gamma \) in the simplex graph \( \mathcal{G}_k \) where \( n(\gamma) \) is the length of \( \gamma \).
Pecha-Kucha is a presentation form in which 20 slides are shown exactly 20 seconds each. Here are slides for a talk given on March 6, 2013. The bold text had been prepared to be spoken, the text after consists of remarks posted on March 2013.

What if Archimedes would have known the concept of a function? The following story borrows from classes taught to students in singular variable calculus and at the extension school. The mathematics is not new but illustrates part of what one calls “quantum calculus” or “calculus without limits”.

Quantum calculus comes in different flavors. A book with that title was written by Kac and Cheung [32]. The field has connections with numerical analysis, combinatorics and even number theory. The calculus of finite differences was formalized by George Boole but is much older. Even Archimedes and Euler thought about calculus that way. It also has connections with nonstandard analysis. In ”Math 1a, a single variable calculus course or Math E320, a Harvard extension school course, I spend 2 hours each semester with such material. Also, for the last 10 years, a small slide show at the end of the semester in multivariable calculus featured a presentation ”Beyond calculus”.

About 20 thousand years ago, mathematicians represented numbers as marks on bones. I want you to see this as the constant function 1. Summing up the function gives the concept of number. The number 4 for example is 1 + 1 + 1 + 1. Taking differences brings us back: 4 – 3 = 1.

The Ishango bone displayed in the slide is now believed to be about 20’000 years old. The number systems have appeared in the next couple of thousand years, independently in different places. They were definitely in place 4000 BC because we have Carbon dated Clay tablets featuring numbers. Its fun to make your own Clay tablets: either on clay or on chewing gum.
When summing up the previous function \( f(x) = x \), we get triangular numbers. They represent the area of triangles. Gauss got a formula for them as a school kid. Summing them up gives tetrahedral numbers, which represent volumes. Integration is summation, differentiation is taking differences.

The numbers \([x]^n/n!\) are examples of **polytopic numbers**, numbers which represent patterns. Examples are triangular, tetrahedral or pentatopic numbers. Since we use the forward difference \( f(x+1) - f(x) \) and start summing from 0, our formulas are shifted. For example, the triangular numbers are traditionally labeled \( n(n+1)/2 \), while we write \( n(n-1)/2 \). We can think about the new functions \([x]^n\) as a new basis in the linear space of polynomials.

The polynomials appear as part of the Pascal triangle. We see that the summation process is the inverse of the difference process. We use the notation \([x]^n/n!\) for the function in the \( n\)th row.

The “renaming idea” \([x]^n\) is part of **quantum calculus**. It is more natural when formulated in a non-commutative algebra, a crossed product of the commutative algebra we are familiar with. The commutative \( X = C^\ast\)-algebra \( C(R) \) of all continuous real functions encodes the topology of the real numbers. If \( s(x) = x + 1 \) is translation, we can look at the algebra of operators on \( H = L^2(R) \) generated by \( X \) and the translation operator \( s \). With \([x] = xs^\ast\), the operators \([x]^n\) are the multiplication operators generated by the polynomials of the slide. The derivative \( Df \) be written as the commutator \( Df = [s,f] \). In quantum mechanics, functions are treated as operators. The deformed algebra is no more commutative: with \( Qf(x) = xs^\ast \) satisfying \( Q^n = [x]^n \) and \( Pf = iDf \) the anti-commutation relations \([Q,P] = QP - PQ = i\) hold. Hence the name “quantum calculus”.
With the difference operation $D$ and the summation operation $S$, we already get an important formula. It tells that the derivative of $x^n$ is $n$ times $x$ to the $n-1$. Let's also introduce the function $\exp(ax) = (1 + a)^x$ to the power $x$. This is the compound interest formula. We check that its derivative is a constant times the function itself.

The deformed exponential is just a rescaled exponential with base $(1 + a)$. Writing it as a compound interest formula allows to see the formula $\exp'(ax) = a \exp(ax)$ as the property that the bank pays you a multiple $a \exp(ax)$ to your fortune $\exp(ax)$. The compound interest formula appears in the movie "The bank". We will below look at the more general exponential $\exp_h(x) = (1 + ah)^{x/h}$ which is the exponential $\exp(ax)$ to the "Planck constant" $h$. If $h$ goes to zero, we are lead to the standard exponential function $\exp(x)$. One can establish the limit because $\exp_h(ax) \leq \exp(ax) \leq (1 + ah) \exp_h(ax)$.

The heart of calculus is the fundamental theorem of calculus. If we take the sum of the differences, we get the difference of $f(x)$ minus $f(0)$. If we take the difference of the sum, we get the function $f(x)$. The pictures pretty much prove this without words. These formulas are true for any function. They do not even have to be continuous.

Even so the notation $D$ is simpler than $d/dx$ and $S$ is simpler than the integral sign, it is the language with $D, S$ which makes the theorem look unfamiliar. Notation is very important in mathematics. Unfamiliar notation can lead to rejection. One of the main points to be made here is that we do not have to change the language. We can leave all calculus books as they are and just look at the content in new eyes. The talk is advertisement for the calculus we know, teach and cherish. It is a much richer theory than anticipated.
The proof of part I shows a telescopic sum. The cancellations are at the core of all fundamental theorems. They appear also in multivariable calculus in differential geometry and beyond.

It is refreshing that one can show the proof of the fundamental theorem of calculus early in a course. Traditionally, it takes much longer until one reaches the point. Students have here only been puzzled about the fact that the result holds only for $x=nh$ and not in general. It would already here mathematically make more sense to talk about a result on a finite linear graph but that would increase mentally the distance to the traditional calculus in [10] on "Historical Reflections on Teaching the fundamental Theorem of Integral Calculus". Bressoud tells that "there is a fundamental problem with the statement of the FTC and that only a few students understand it". I hope this Pecha-Kucha can help to see this important theorem from a different perspective.

The proof of the second part is even simpler. The entire proof can be redone, when the step size $h=1$ is replaced by a positive $h$. The fundamental theorem allows us to solve a difficult problem: summation is global and hard, differences is local and easy. By combining the two we can make the hard stuff easy.

This is the main point of calculus. Integration is difficult, Differentiation is easy. Having them linked makes the hard things easy. At least part of it. It is the main reason why calculus is so effective. These ideas go over to the discrete. Seeing the fundamental idea first in a discrete setup can help to take the limit when $h$ goes to zero.
We can adapt the step size $h$ from 1 to become any positive number $h$. The traditional fundamental theorem of calculus is stated here also in the notation of Leibniz. In classical calculus we would take the limit $h$ to 0, but we do not go that way in this talk.

When teaching this, students feel a slight unease with the additional variable $h$. For mathematicians it is natural to have free parameters, for students who have had little exposure to mathematics, it is difficult to distinguish between the variable $x$ and $h$. This is the main reason why in this talk, I mostly stuck to the case $h=1$. In some sense, the limit $h$ to 0 is an idealization. This is why the limit $h$ going to 0 is so desirable. But we pay a prize: the class of functions we can deal with, is much smaller. I myself like to see traditional calculus as a limiting case of a larger theory. The limit $h$ to zero is elegant and clean. But it takes a considerable effort at first to learn what a limit is. Everybody who teaches the subject can confirm the principle that often things which were historically hard to find, are also harder to teach and master. Nature might have given up on it early on: $10^{-37}$ seconds after the big bang: “Darn, lets just do it without limits ...!”

We can define cosine and sine by using the exponential, where the interest rate is the square root of $-1$. These deformed trigonometric functions have the property that the discrete derivatives satisfy the familiar formula as in the books. These functions are close to the trigonometric functions we know if $h$ is small.

This is by far the most elegant way to introduce trigonometric functions also in the continuum. Unfortunately, due to lack of exposure to complex numbers, it is rarely done. The Planck constant is $h = 6.626068 \cdot 10^{-34} m^2 kg/s$. The new functions $(1 + iha)^{x/h} = \cos(a \cdot x) + i \sin(a \cdot x)$ are then very close to the traditional cos and sin functions. The functions are not periodic but a growth of amplitude is only seen if $xa$ is of the order $1/h$. It needs even for X-rays astronomical travel distances to see the differences.
The fundamental theorem of calculus is fantastic because it allows us to sum things which we could not sum before. Here is an example: since we know how to sum up the deformed polynomials, we can find formulas for the sum of the old squares. We just have to write the old squares in terms of the new squares which we know how to integrate.

This leads to a bunch of interesting exercises. For example, because \( x^2 = [x^2] + [x] \), we have \( Sx^2 = S[x^2] + S[x] = [x^3]/3 + [x^2]/2 \) we get a formula for the sum of the first \( n-1 \) squares. Again we have to recall that we sum from 0 to \( n-1 \) and not from 1 to \( n \). By the way: this had been a 'back and forth' in the early lesson planning. I had started summing from 1 to \( n \) and using the backwards difference \( Df(x) = f(x) - f(x - 1) \). The main reason to stick to the forward difference and so to polynomials like \( [x]^9 = x(x-h)(x-2h)...(x-8h) \) was that we are more familiar with the difference quotient \([f(x+h)-f(x)]/h\) and also with left Riemann sums. The transition to the traditional calculus becomes easier because difference quotients and Riemann sums are usually written that way.

The integration rule analogue to integration by parts is called **Abel summation** which is important when studying Dirichlet series. Abel summation is used in this project. The Leibniz formula has a slight asymmetry. The expansion of the rectangle has two main effects \( fDg \) and \( Dfg \), but there is an additional small part \( DfDg \). This is why we have \( D(fg) = fDg + Dfg \). The formula becomes more natural when working in the non-commutative algebra mentioned before: if \( Df = [s, f] \), then \( D(fg) = fDg + Dfg \) because the translation operator \( s \) now takes care of the additional shift: \( (Df)g = (f(x+1) - f(x))g(x+1) \). The algebra picture also explains why \([x]^n[x]^m\) is not \([x]^{n+m}\) because also the multiplication operation is deformed in that algebra. In the noncommutative...
algebra we have \((xs)^n(xs)^m = (xs)^{n+m}\). While the algebra deformation is natural for mathematicians, it can not be used in calculus courses. This can be a reason why it appears strange at first, like quantum mechanics in general.

The chain rule also looks the same. The formula is exact and holds for all functions. Writing down this formula convinces you that the chain rule is correct. The only thing which needs to be done in a traditional course is to take the limit.

Most calculus textbooks prove the chain rule using linearization: first verify it for linear functions then argue that the linear part dominates in the limit. [Update of December 2013: The proof with the limit.] Reversing the chain rule leads to the integration tool substitution. Substitution is more tricky here because as we see different step sizes \(h\) and \(H\) in the chain rule. This discretization issue looks more serious than it is. First of all, compositions of functions like \(\sin(\sin(x))\) are remarkably sparse in physics. Of course, we encounter functions like \(\sin(kx)\) but they should be seen as fundamental functions and defined like \(\sin(k \cdot x) = \text{Im}(1 + khi)^{x/h}\). This is by the way different from \(\sin(kx) = (1 + khi)^{(kx)/h}\).

Also the Taylor theorem remains true. The formula has been discovered by Newton and Gregory. Written like this, using the deformed functions, it is the familiar formula we know. We can expand any function, continuous or not, when knowing the derivatives at a point. It is a fantastic result. You see the start of the proof. It is a nice exercise to make the induction step.

**James Gregory** was a Scottish mathematician who was born in 1638 and died early in 1675. He has become immortal in the Newton-Gregory interpolation formula. Despite the fact that he gave the first proof of the fundamental theorem of calculus, he is largely unknown. As a contemporary of Newton, he was also in the shadow of Newton. It is interesting how our
value system has changed. “Concepts” are now considered less important than proofs. Gregory would be a star mathematician today. The Taylor theorem can be proven by induction. As in the continuum, the discrete Taylor theorem can also be proven with PDE methods:

\[ D_t f = Df, \]

where \( D_t f(x, t) = f(x, t + 1) - f(x, t) \) so that

\[ f(0, t) = \left[ \exp(Dt)f \right](0) = f(0) + Df(0)/1! + D^2f(0)/2! + \ldots. \]

Here is an example, where we expand the exponential function and write it as as a sum of powers. Of course, all the functions are deformed functions. The exponential function as well as the polynomials \( x^n \) were newly defined. We see what the formula means for \( x = 5 \).

It is interesting that Taylor’s theorem has such an arithmetic incarnation. Usually, in numerical analysis text, this Newton-Gregory result is treated with undeformed functions, which looks more complicated. The example on this slide is of course just a basic property of the Pascal triangle. The identity given in the second line can be rewritten as 32 = 1 + 5 + 10 + 10 + 5 + 1 which in general tells that if we sum the n’th row of the Pascal triangle, we get \( 2^n \). Combinatorially, this means that the set of all subsets can be counted by grouping sets of fixed cardinality. It is amusing to see this as a Taylor formula (or Feynman-Kac). But it is more than that: it illustrates again an important point I wanted to make: we can not truly appreciate combinatorics if we do not know calculus.

Taylor’s theorem is useful for data interpolation. We can avoid linear algebra and directly write down a polynomial which fits the data. I fitted here the Dow Jones data of the last 20 years with a Taylor polynomial. This is faster than data fitting using linear algebra.

When fitting the data, the interpolating function starts to oscillate a lot near the ends. But that’s not the deficit of the method. If we fit n data points with a polynomial of degree n-1, then we get a unique solution. It is the same solution as the Taylor theorem gives. The proof of the theorem is a simple induction
step, using a property of the Pascal triangle. Assume, we know \( f(x) \). Now apply \( D \), to get \( Df(x) \). We can now get \( f(x+1) = f(x) + Df(x) \). Adding the terms requires a property of the Pascal triangle.

We can also solve differential equations. We can use the same formulas as you see in books. We deform the operators and functions so that everything stays the same. Difference equations can be solved with the same formulas as differential equations.

To make the theory stronger, we need also to deform the log, as well as rational functions. This is possible in a way so that all the formulas we know in classical calculus holds: first define the log as the inverse of \( \exp \). Then define \( [x]^{-1} = D \log(x) \) and \( [x]^{-n} = D[x]^{1-n}/(1-n) \). Now \( Dx^n = nx^{n-1} \) holds for all integers. We can continue like that and define \( \sqrt{x} \) as the inverse of \( x^2 \) and then the \( x^{-1/2} \) as \( 2D \sqrt{x} \). It is calculus which holds everything together.

In multivariable calculus, space takes the form of a graph. Scalar functions are functions on vertices, Vector fields are signed functions on oriented edges. The gradient of a function is a vector field which is the difference of function values. Integrating the gradient of a function along a closed path gives the difference between the potential values at the end points. This is the fundamental theorem of line integrals.

This result is also in the continuum the easiest version of Stokes theorem. Technically, one should talk about 1-forms instead of vector fields. The one forms are anti-commutative functions on edges. \([40]\) is an exhibit of three theorems (Green-Stokes, Gauss-Bonnet \([37]\), Poincaré-Hopf \([38]\)), where everything is defined and proven on two pages. See also the overview \([42]\).
Stokes theorem holds for a graph for which the boundary is a graph too. Here is an example of a "surface", which is a union of triangles. The curl of a vector field $F$ is a function on triangles is defined as the sum of vector fields along the boundary. Since the terms on edges in the intersection of triangles cancel, only the line integral along the boundary survives.

This result is old. It certainly seems have been known to Kirchhoff in 1850. Discrete versions of Stokes pop up again and again over time. It must have been Poincaré who first fully understood the Stokes theorem in all dimensions and in the discrete, when developing algebraic topology. He introduced chains because unlike graphs, chains are closed under boundary operation. This is a major reason, algebraic topologists use it even so graphs are more intuitive. One can for any graph define a discrete notion of differential form as well as an exterior derivative. The boundary of graph is in general only a chain. For geometric graphs like surfaces made up of triangles, the boundary is a union of closed paths and Stokes theorem looks like the Stokes theorem we teach.

Could Archimedes have discovered the fundamental theorem of calculus? His intellectual achievements surpass what we have seen here by far. Yes, he could have done it. If he would not have been given the sword, but the concept of a function.

Both the precise concept of limit as well as the concept of functions had been missing. While the concept of limit is more subtle, the concept of function is easier. The basic ideas of calculus can be explained without limits. Since the ideas of calculus go over so nicely to the discrete, I believe that calculus is an important subject to teach. It is not only a benchmark and a prototype theory, a computer scientist who has a solid understanding of calculus or even differential topology can also work much better in the discrete. To walk the talk: source code of a program in computer vision was written from scratch, takes a movie and finds and tracks features like corners [43]. There is a lot of calculus
used inside.

This was a talk given on March 6 at a “Pecha-Kucha” event at the Harvard Mathematics department, organized by Sarah Koch and Curt McMullen. Thanks to both for running that interesting experiment.

4. Problems in single variable calculus

Teachers at the time of Archimedes could have assigned problems as follows. Some of the following problems have been used in single variable and extension school classes but most are not field tested. It should become obvious however that one could build a calculus course sequence on discrete notions. Since the topic is not in books, it would be perfect also for an “inquiry based course”, where students develop and find theorems themselves.

1. Compute $Dg(n)$ with $g(n) = [n]^4 = (n - 3)(n - 2)(n - 1)n$ and verify that $Dg(n) = 4[n]^3$ with $[n]^3 = (n - 2)(n - 1)n$.

2. Find $Sf(n)$ with $f(n) = n$. This was the task, which the seven year old Gauss was given by his school teacher Herr Büttner. Gauss computed $Sf(101) = 5050$ because he saw that one can pair 50 numbers summing up to 101 each and get $[101]^2/2 = 101 * 100/2$. Legend tells that he wrote this number onto his tablet and threw it onto the desk of his teacher with the words “Hier ligget sie!” (=“Here it is!”). Now its your turn. You are given the task to sum up all the squares from 1 to 1000.

3. Find the next term in the sequence

2, 10, 30, 68, 130, 222, 350, 520, 738, 1010, 1342,...

by taking “derivatives” until you see a pattern, then “integrate” that pattern.

4. Lets compute some trig function values. We know for example. $\sin(0) = 0, \sin(1) = 1, \sin(2) = 2, \sin(3) = 1$ and $\cos(0) = 1, \cos(1) = 1, \cos(2) = 0, \cos(3) = -1$. Compute $\sin(10)$ and $\cos(10)$. (Answer: 32 and 0.)

5. Find $F(n) = Sf(n)$ for the function $f(n) = n^2$. Your explicit formula should satisfy $F(1) = 1, F(2) = 5, F(3) = 14$ and lead to the sequence

0, 1, 5, 14, 30, 55, 91, 140, 204, 285,...

6. Find $g(x) = Df(x)$ for $f(x) = 3[x]^5 + 3^x - 2x + 7$ and evaluate it at $x = 10$. Answer: We have seen that $D(1 + a)^x = a(1 + a)^x$ so that $D3^x = 2 \cdot 3^x$. $g(x) = 15[x]^4 + 2 \cdot 3^x - 2$. We have $g(10) = 15(10*9*8*7) + 23^{10} - 2 = 193696$. Indeed:
Verify that \( \frac{d}{dx} f = f \) for \( f = \exp(x) \) and conclude that \( S f(x) = f(x + 1) - 1 \). For example, for \( x = 4 \) we have \( 1 + 2 + 4 + 8 = 16 - 1 \).

The Fibonacci sequence \( 1, 1, 2, 3, 5, 8, 13, 21, \ldots \) satisfies the rule \( f(x) = f(x - 1) + f(x - 2) \). Verify that \( \frac{d}{dx} f(x) = f(x - 1) \) and use this fact to show that \( S f(x) = f(x + 1) - 1 \). For example, for \( x = 4 \) we get \( 1 + 1 + 2 + 3 = 8 - 1 \) or for \( x = 6 \) we have \( 1 + 1 + 2 + 3 + 5 + 8 = 21 - 1 \).

Verify that the Taylor series \( \exp(x) = 1 + x + x^2/2! + x^3/3! + \ldots \) sums up the \( n \)'th line in the Pascal triangle. For example, \( \exp(4) = 1 + 4 + 6 + 4 + 1 = 16 \).

To do so, check that \( \frac{d^n}{dx^n} f \) is only since the Planck constant equal to \( 1 \). Let's look at the harmonic oscillator with \( V = \frac{1}{2} k x^2 \). The simplest way to derive the double angle formulas \( \cos(2x) = \cos^2(x) - \sin^2(x) \) and \( \sin(2x) = 2 \cos(x) \sin(x) \). Check whether these double angle formulas work also with the new functions \( \cos(x), \sin(x) \). Be careful that \( \cos(a \cdot x) = \text{Re}(1 + ia)^x \) is not the same than \( \cos(ax) = \text{Re}(1 + i)^{ax} \).

Verify that \( \sin(x) \) has roots at \( 4k \) with \( k \in \mathbb{Z} \) and \( \cos(x) \) has roots at \( 2 + 4k \) with \( k \in \mathbb{Z} \). To do so, note \( \sin(x) = \text{Im}(e^{x\log(1+i)}) = \text{Im}(e^{x(\sqrt{2}+i\pi/4)}) \).

Find the solution of the harmonic oscillator \( \frac{d^2}{dx^2} f = -9f \) which has the initial conditions \( f(0) = 2, f(1) = 5 \). Solution: the solution is of the form \( a \cos(3x) + b \sin(3x) \), where \( a, b \) are constants. Taking \( x = 0 \) gives \( a = 2 \). Taking \( x = 1 \) gives \( 2 \cos(3) + b \sin(3) = 5 \). Since \( \cos(3) = -2 \) and \( \sin(3) = 2 \) we have \( b = 9/2 \). The solution is \( f(x) = 2 \cos(3x) + 9 \sin(3x)/2 \).

If \( h \) is a small positive number called Planck constant. Denote by \( \sin, \cos \) the classical sin and cos functions. Verify that \( |\sin(h \cdot (x/h)) - \sin(x)| \leq xh \).

Solution: \( |\sin((1+ih)x/h - \sin(x)| = |\sin(x + ih)\log(1+ih) - \sin(x)| \). Remark. With \( h = 7 \cdot 10^{-34} \), a Gamma ray of wave length \( 10^{-12} \) traveling for one million light years \( 10^6 \cdot 10^{16} \) meters, we start to see a difference between \( \sin(h \cdot x/h) \) and \( \sin(x) \). The amplitude of the former will grow and be seen as a gamma ray burst.

The functions \( \sin(x) \) and \( \cos(x) \) are not periodic on the integers. While we know \( \sin(4k) = 0 \) for every integer \( k \), we have \( \sin(4k + 1) \) exploding exponentially (this is only since the Planck constant equal to 1). Let's look at the tangent function.
\[ \tan(x) = \frac{\sin(x)}{\cos(x)}. \]
Verify that it is 4 periodic \(0, 1, \infty, -1, 0, 1, \infty, -1, \ldots\) on the integers.

16 Let's call \(x\) be a **discrete gradient critical point** of \(f\) if \(Df(x) = 0\). Assume that \(f\) is a continuous function from the real line \(R\) to itself and that has a discrete gradient critical point \(a\). Verify that there is a classical local maximum or classical local minimum of \(f\) in \((a, a + 1)\). Solution: this is a reformulation of the classical **Bolzano extremal value theorem** on the interval \([a, a + 1]\).

17 Injective functions on \(Z\) do not have discrete gradient critical points. Critical points must be defined therefore differently if one only looks at functions on a graph. A point \(x\) is a **critical point** if the index \(i_f(x) = 1 - |S(x) \cap \{f(y) < f(x)\}|\) is not zero. In other words, if either both or no neighbor of \(x\) have smaller values than \(f(x)\) then \(x\) is a critical point. Verify that if \(f\) is an injective \(n\) periodic function on \(Z\), then \(\sum_{k=1}^{n} i_f(k)\) is equal to 0. **Hint:** use induction. Assume it is true for \(n\), look at \(n + 1\), then see what happens if the maximum \(x\) of \(f\) is taken away.

18 If \(f\) is a classically differentiable function, then for any \(a\), there exists \(x \in [a, a + 1]\) such that \(f'(x) = Df(a)\). Here \(f'(x)\) is the classical derivative and \(Df(x) = f(x + 1) - f(x)\) the discrete derivative. Solution: this is a reformulation of the classical **mean value theorem**.

19 The **Taylor expansion** \(p(x) = \sum_k D^k f(0) X^k / k!\) extends a function \(f\) of compact support from the integers to the real line. It is an explicit **interpolation formula** which provides us with a function \(f(x)\) on the real line satisfying \(p(x) = f(x)\) for integer arguments \(x\). Perform the Taylor expansion of the function \(f(x) = x^2\).

20 Check that unlike the classical trigonometric function \(\sin(x)\) the function \(\sin(x)\) is not **odd** and the function \(\cos(x)\) is not **even**. Indeed, check that \(\sin(x) \to 0\) for \(x \to -\infty\) while \(\sin(x)\) becomes unbounded for \(x \to \infty\).

21 Compute the sum \(\sum_{x=0}^{9} \sin(3 \cdot x)\). **Answer:** \(-33237 = (1 - \cos(3 \cdot 10))/3\). Now give an analytic expression for \(\sum_{k=0}^{100} \sin(3 \cdot x)\) which can no more be evaluated by computing the sum directly.

22 Let us in this problem also write \(f' = Df\). Integrating the **product rule** \(D(fg) = Df g + f Dg\) gives \(fg = S(f'g) + S(f g')\) leads to **integration by parts** formula

\[ S(f'g) = fg - S(f^+ g'), \]
an expression which is in the discrete also called **Abel summation**. It is usually written as

\[
\sum_{k=0}^{n-1} f'_k g_k = (f_n g_n - f_0 g_0) - \sum_{k=0}^{n-1} f_{k+1} g'_k.
\]

Use integration by parts to sum up \(\sum_{x=0}^{102} \sin(x) x\). **Answer**: the sum is \((\sin(103) - \sin(1)) - 102 \cos(103) = -231935380809580545\).

23 We have defined \(1/X\) as the derivative of \(\log(x)\). Verify that \(1/X = \log(1 + 1/x)/\log(2)\), where \(\log\) is the good old classical natural logarithm.

### 5. Problems in multi variable calculus

1 **Bolzano's extremal value theorem** is obvious for finite graphs because a function on a finite set always has a global maximum and minimum. Verify that for a global minimum, the **index** \(i_f(x)\) is equal to 1. For every integer \(n\) find a graph \(G\) and a function \(f\) for which a maximum is equal to \(n\). For example, for a star graph with \(n\) rays and a function \(f\) having the maximum in the center \(x\), we have \(i_f(x) = 1 - n\).

2 A **curve** in \(G\) is a subgraph \(H\) with the property that every vertex has either 1 or 2 neighbors. The **arc length** of a path is the number of edges in \(H\). The boundary points are the vertices with one neighbor. A connected **surface** \(H\) in \(G\) is a subgraph for which every vertex has either a linear graph or circular graph as a unit sphere. The surface area of \(H\) is the number of triangles in \(H\). The interior points are the points for which the unit sphere is a circular graph. A surface is called **flat** if every interior point has a unit sphere which is \(C_6\). Flat surfaces play the role of regions in the plane. The curvature of a boundary point of a flat surface is \(1 - |S(x)|/6\). Verify that the sum of the curvatures of boundary points is \(1 - g\) where \(g\) is the number of holes. This is a discrete **Hopf Umlaufsatz**.

3 A graph \(G\) is called **simply connected** if every simple and connected curve \(C\) in \(G\) can be deformed to a single edge using the following two **deformation steps** of a \(C\): replace two edges of \(C\) within a triangle with the third edge of that triangle or replace an edge of \(C\) in a triangle with the two other edges of the triangle. A deformation step is called **valid** if the curve remains a simple connected curve or degenerates to \(K_3, K_2\) after the deformation. A 1-form \(F\) is called a **gradient field**, if \(F = df\) for some scalar function \(f\). Verify that on a simply connected surface, a 1-form \(F\) is a gradient field if and only if the curl \(dF\) is zero on each triangle.

4 **Stokes theorem** for a flat orientable surface \(G\) is what traditionally is called **Greens theorem**. It tells that the sum of the curls over the triangles is equal
to the line integral along the boundary. Look at the wheel graph \( W_6 \) which is obtained by adding an other vertex \( x \) to a circular graph \( C_6 \) and connecting all vertices \( y_i \) of \( C_6 \) with \( x \). Define the 1-form \( F \), where each of the central edges \((x,y_i)\) is given the value 1 and each edge \((y_i,y_{i+1})\) is given the value \( i \). Find the curl \( dF \) on all the triangles and verify Green’s theorem in that case.

5. A graph \( G \) is a surface if for every vertex \( x \in G \), the unit sphere \( S(x) \) is a circular graph. The curvature at a point \( x \) of a surface is defined as \( 1 - |S(x)|/6 \). Verify that an octahedron is a discrete surface and that the curvature is constant \( 1/3 \) everywhere. See that the curvatures add up to 2. Verify also that the icosahedron is a discrete surface and that the curvature is constant \( 1/6 \) adding up to 2. You have verified two special cases of Gauss-Bonnet. Play with other discrete surfaces but make sure that all “faces” are triangles.

6. We want to see an analogue of the second derivative test for functions \( f(x,y) \) of two variables on a discrete surface \( G \) like an icosahedron. For an injective function \( f \) on the vertices, call \( x \) a critical point with nonzero index if \( S^-(x) = S(x) \cap \{ y \mid f(y) < f(x) \} \) has Euler characteristic different from 1. The number \( i_f(x) = 1 - \chi(S^-(x)) \) plays the role of the discriminant. Verify that if it is negative then we deal with a saddle point which is neither maximum nor minimum. Verify that if it is positive with empty \( S^-(x) \), then it is a local minimum. Verify that if the discriminant is positive and \( S^-(x) \) is a circular graph, then it is a local maximum. As in the continuum, where the discriminant at critical points can be zero, this can happen also here but only in higher dimensions. Maybe you can find an example of a function on a three dimensional graph.

7. Given a surface \( G = (V,E) \) and an injective function \( f \) on \( V \), define the level curve \( C \) of \( f = c \) as the graph which has as vertices the edges of \( G \) for which \( f \) changes. Two such vertices are connected if they are contained in a common triangle. (Graph theorists call this a subgraph of the “line graph” of \( G \).) Verify that for a two dimensional geometric graph, level curve \( C \) of \( f \) is a finite union of circular graphs and so a geometric one dimensional graph. The fact that the gradient \( df(x) \) is a function on \( C \) changing sign on every original edge is the analogue of the classical fact that the gradient is perpendicular to the level curve.

8. An oriented edge \( e \) of a graph \( G \) is a "direction" and plays the role of unit vectors in classical calculus. The value of \( df(e) \) is the analogue of a directional derivative. Start with a vertex \( x \) of \( G \), then go into direction \( e \), where the directional derivative \( df(e) \) is maximal and stop if there is no direction in which the directional derivative is positive. This discrete gradient flow on the vertex set leads to a local maximum. Investigate how this notion of directional derivative corresponds to the notion \( D_v f(x) = \nabla f(x) \cdot v \) in the continuum. Especially note that for an injective function, the directional derivative is never zero.
An injective function \( f \) on a graph has no points for which \( df(a) = 0 \). We therefore have to define critical points differently. We say \( x \) is a critical point of \( f \) if \( i_f(x) = 1 - \chi(S_f(x)) \) is not zero. Let \( G \) be a discretisation of a disc. Check that for maxima, \( i_f(x) = 1 \) and \( S_f(x) \) is a circular graph. Check that for minima, \( i_f(x) = 1 \) and \( S_f(x) \) is empty. For saddle points, \( i_f(x) = -1 \) where \( S_f(x) \) consists of two linear graphs. A vertex is called a monkey saddle if \( i_f(x) = -2 \). Verify the "island theorem": the sum over all indices over the disc is equal to 1. A discrete Poincaré-Hopf theorem assures that in general \( \sum_x i_f(x) = \chi(G) \).

Let \( F,G \) be two 1-forms, define the dot product of \( F,G \) at a vertex \( x \) as \( \langle F,G \rangle(x) = \sum_{e \in E(x)} F(e)G(e) \), where \( E(x) \subset E \) is the set of edges attached to \( x \). Define the length \( |F| \) as \( \sqrt{\langle F,F \rangle} \). The angle between two 1-forms at \( x \) is defined by the identity \( \langle F,G \rangle = |F| \cdot |G| \cos(\alpha) \). Given two functions \( f,g \) and an edge \( e = (a,b) \) define the angle between two hyper surfaces \( f = c, g = d \) at an edge \( e \) as the angle between \( df(e) \) and \( dg(e) \). The hyper surface \( f = c \) is part of the line graph of \( G \): it has as vertices the edges of the graph at which \( f \) changes sign and as edges the pairs of new vertices which intersect when seen as old edges.

Let \( F,G \) be two 1-forms. Define the cross product of \( F,G \) at a vertex \( x \) as the function on adjacent triangle \( (x,y,z) \) as \( F \times G(x,y,z) = F(x,y)G(x,z) - F(x,z)G(x,y) \). It is a number which is the analogue of projecting the classical cross product onto a 2-dimensional direction in space. The cross product is anti-commutative. Unlike a two form, it is not a function on triangles because the order of the product plays a role. It can be made a 2-form by antisymmetrization. Here is the problem: verify that for any two functions \( f,g \) on the vertices of a triangle the cross product \( df \times dg \) is up to a sign at all vertices.

Given three 1-forms \( F,G,H \), define the triple scalar product at a vertex \( x \) as a function on an adjacent tetrahedron \( (x,y,z,w) \) as \( F \cdot G \times H(x,y,z,w) = \det \begin{bmatrix} F(x,y) & F(x,z) & F(x,w) \\ G(x,y) & G(x,z) & G(x,w) \\ H(x,y) & H(x,z) & H(x,w) \end{bmatrix} \). Verify that this triple scalar product has the same properties as the triple scalar product in classical multivariable calculus.

Let \( G = K_2 \) be the complete graph with two vertices and one edge. Verify that the Dirac operator \( D \) and form Laplacian \( L = D^2 \) are given by

\[
D = \begin{bmatrix} 0 & 0 & -1 \\ 0 & 0 & 1 \\ -1 & 1 & 0 \end{bmatrix}, \quad L = \begin{bmatrix} 1 & -1 & 0 \\ -1 & 1 & 0 \\ 0 & 0 & 2 \end{bmatrix}.
\]
For an informal introduction to the Dirac operator on graphs and Mathematica code, see [31]. The Laplacian has the eigenvalues $2, 2, 0$. Write down the solution of the wave equation $f'' = -Lf$ where $f(t) = (u(t), v(t), w(t))$ and $u(t), v(t)$ are the values of $f$ on vertices and $w(t)$ is the evolution of the 1-form.

14 The graph $G = K_5$ is the smallest possible four dimensional space. Given a 1-form $A$, we have at each vertex 6 triangles. The 2-form $F = dA$ defines so a field with 6 components. If one of the four edges emerging from $x$ is called "time" and the 3 triangles containing this edge define the electric field $E_1, E_2, E_3$ and the other 3 triangles are the magnetic field $B_1, B_2, B_3$. The equations $dF = 0, d^*F = j$ are called Maxwell equations. Can we always get $F$ from the current $j$? Verify that if $\text{div}(A) = d^*A = 0$ (Coulomb gauge), then the Poisson equation $LA = j$ always allows to compute $A$ and so $F = dA$. The Coulomb gauge assumption is one of Kirchhoff’s conservation law: the sum over all in and outgoing currents at each node are zero.

15 Let $G = C_5$ denote the circular graph with 5 vertices. Find a function on the vertex set, which has exactly two critical point, one minimum of index 1 and one maximum of index $-1$. Now find a function with 4 critical points and verify that there are no functions with an odd number of critical points. **Hint:** show first that the index of a critical point is either 1 or $-1$ and use then that the indices have to add up to 0.

16 Now look at the tetrahedron $G = K_5$. Verify that every function on $G$ has exactly one critical point, the minimum.

17 Take a linear graph $L_3$ which consists of three vertices and two edges. There are 6 functions which take values $\{1, 2, 3\}$ on the vertices. In each case compute the index $i_f(x)$ and verify that the average over all these indices gives the curvature of the graph, which is $1/2$ at the boundary and 0 in the center. This is a special case of a result proven in [39].

**Appendix A: Graph theory glossary**

A **graph** consists of a vertex set $V$ and an edge set $E \subset V \times V$. If $V$ and $E$ are finite, it is called a **finite graph**. If every $e = (a, b) \in E$ satisfies $a \neq b$, elements $(a, b)$ and $(b, a)$ are identified and every $e$ appears only once, the graph is called a **finite simple graph**. In other words, we disregard loops $(a, a)$, look at undirected graphs without multiple connections. The cardinality of $V$ is called **order** and the cardinality of $E$ is called **size** of the graph. When counting $E$, we look at the unordered pairs $e = \{a, b\}$ and not the ordered pairs $(a, b)$. The later would be twice as large.

Restricting to finite simple graphs is not much of a restriction of generality. The notion of loops and multiple connections can be incorporated later by introducing the algebraic notion of **chains** defined over the graph, directions using
Given a finite simple graph $G = (V, E)$, the elements in $V$ are called vertices or nodes while the elements in $E$ are called edges or links. A finite simple graph contains more structure without further input. A graph $H = (W, F)$ is called a subgraph of $G$ if $W \subseteq V$ and $F \subseteq E$. A graph $(V, E)$ is called complete if the size is $n(n-1)/2$, where $n$ is the order. In other words, in a complete graph, all vertices are connected. The dimension of a complete graph is $n-1$ if $n$ is the order. A single vertex is a complete graph of dimension 0, an edge is a complete graph of dimension 1 and a triangle is a complete graph of dimension 2. Let $G_k$ denote the set of complete subgraphs of $G$ of dimension $k$. We have $G_0 = V, G_1 = E$ and $G_2$ the set of triangles in $G$ and $G_3$ the set of tetrahedra in $G$. Let $v_k$ be the number of elements in $G_k$. The number $\chi(G) = \sum_{k=0}^{\infty} (-1)^k v_k$ is called the Euler characteristic of $G$. For a graph without tetrahedra for example, we have $\chi(G) = v - e + f$, where we wrote $v = v_0, e = v_1, f = v_2$. For example, for an octahedron or icosahedron, we have $\chi(G) = 2$. More generally, any triangularization of a two dimensional sphere has $\chi(G) = 2$. The cube graph has no triangles so that $\chi(G) = 8 - 12 = -4$ illustrating that we see it as a sphere with $6 = 2g$ holes punched into. For a discretisation of a sphere with $2g$ holes, we have $\chi(G) = 2 - 2g$. If $g = 1$ for example, where we have 2 holes, we can identify the boundaries of the holes without changing the Euler characteristic and get a torus.

Often, graphs are considered as "curves", one dimensional objects in which higher dimensional structures are neglected. The most common point of view is to disregard even the two-dimensional structures obtained from triangles and define the Euler characteristic of a graph as $\chi(G) = v - e$, where $v$ is the number of vertices and $e$ the number of edges. For a connected graph this is then often written as $1 - g$, where $g$ is called the genus. It is a linear algebra computation to show that the Euler characteristic $\sum_{k=0}^{\infty} (-1)^k v_k$ is equal to the cohomological Euler characteristic $\sum_{k=0}^{\infty} (-1)^k b_k$, where $b_k$ is the dimension of the cohomology group $H^k(G)$. For example, $b_0$ is the number of connected components of the graph and $b_1$ is the "number of holes" the genus of the curve. Ignoring higher dimensional structure also means to ignore higher dimensional cohomologies which means that the Euler-Poincaré formula is $v_0 - v_1 = b_0 - b_1$ or $v - e = 1 - g$.

Given a subset $W$ of the vertex set, it generates a graph $H = (W, F)$ which has as edge set $F$ all the pairs $(a, b)$ of $W$ with $(a, b) \in E$. A vertex $y$ is called a neighbor of $x$ if $(x, y) \in E$. If $x$ is a vertex of $G$, denote by $S(x)$ the unit sphere of of $x$. It is the graph generated by the set of neighbors of $x$. Given a subgraph $H = (W, F)$ of $G = (V, E)$, we can build a new graph $G = (V \cup \{v\}, E \cup \{(v, y) \mid y \in W\})$ called pyramid extension over $H$. We can define a class of contractible graphs as follows: the one point graph is contractible. Having defined contractibility for graphs of order $n$, a graph $G$ of order $n + 1$ is contractible if there exists a contractible subgraph $H$ of
\(G\) such that \(G\) is a pyramid extension over \(H\). Having a notion of contractibility allows to define critical point properly. And critical points are one of the most important notions in calculus. Given an injective function \(f\) on the vertex set \(V\), a vertex \(x\) is called a critical point of \(f\) if \(S^{-1}(x) = S(x) \cap \{y \mid f(y) < f(x)\}\) is empty or not contractible. Because contractible graphs have Euler characteristic 1, vertices for which \(i_f(x)\) is nonzero are critical points. As in the continuum, there are also critical points with zero index.

Lets look at examples of graphs: the complete graph \(K_{n+1}\) is an \(n\) dimensional simplex. They all have Euler characteristic 1. The circular graph \(C_n\) is a one-dimensional geometric graph. The Euler characteristic is 0. Adding a central vertex to the circular graph and connecting them to all the \(n\) vertices in \(C_n\) produces the wheel graph \(W_n\). If the original edges at the boundary are removed, we get the star graph \(S_n\). All the \(W_n\) and \(S_n\) are contractible and have Euler characteristic 1. The linear graph \(L_n\) consists of \(n + 1\) vertices connected with \(n\) edges. Like the star graph it is also an example of a tree. A tree is a graph without triangles and subgraphs \(C_n\). Trees are contractible and have Euler characteristic 1. The octahedron and icosahedron are examples of two dimensional geometric graphs. Like the sphere, they have Euler characteristic 2.

**Appendix B: Historical notes**

For more on the history of calculus and functions and people, see [11, 7, 9, 18, 8, 17, 33, 5, 58, 6, 62, 16].

Archimedes (287-212 BC) is often considered the father of integral calculus as he introduced Riemann sums with equal spacing. Using comparison and exhaustion methods developed earlier by Eudoxos (408-347 BC), he was able to compute volumes of concrete objects like the hoof, domes, spheres or areas of the circle or regions below a parabola. Even so the concept of function is now very natural to us, it took a long time until it entered the mathematical vocabulary. [8] argues that tabulated function values used in ancient astronomy as well as tables of cube roots and the like can be seen as evidence that some idea of function was present already in antiquity. It was Francois Viètte (1540-1603) who introduced variables leading to what is now called “elementary algebra” and René Descartes (1596-1650) who introduced analytical geometry as well as Pierre de Fermat (1601-1665) who studied also extremal problems and Johannes Kepler (1571-1630) who acceleration as rate of change of velocity and also optimized volumes of wine casks or computed what we would call today elliptic integrals or trigonometric functions.

As pointed out for example in [15], new ideas are rarely without predecessors. The notion of coordinate for example has been anticipated already by Nicole Oresme (1323-1392) [9]. Dieudonné, who praises the idea of function as a “great advance in the
seventeenth century”, points out that first examples of plane curves defined by equations can be traced back to the fourth century BC. While ancient Greek mathematicians already had a notion of tangents, the notion of derivative came only with Isaac Barrow (1630-1703) who computed tangents with a method we would today call implicit differentiation. Isaac Newton introduced derivatives on physical grounds using different terminology like “fluxions”. Until the 18th century, functions were treated as an analytic expression and only with the introduction of set theory due to Georg Cantor, the modern function concept as a single-valued map from a set \( X \) to a set \( Y \) appeared. From [49]: ”The function concept is one of the most fundamental concepts of modern mathematics. It did not arise suddenly. It arose more than two hundred years ago out of the famous debate on the vibrating string”. The word “function” as well as modern calculus notation came to us through Gottfried Leibniz (1646-1716) in 1673. ([8] tells that the term “function” appears the first time in print in 1692 and 1694). Trig functions entered calculus with Newton in 1669 but were included in textbooks only in 1748 [33]. Peter Dirichlet (1805-1859) gave the definition still used today stressing that for the definition \( x \rightarrow f(x) \) it is irrelevant on how the correspondence is established. The notation \( f(x) \) as a value of \( f \) evaluated at a point \( x \) came even later through Leonard Euler (1707-1783) and functions of several variables only appeared at the beginning of the eighteenth century.

When Fourier series became available with Jean Baptiste Fourier (1768-1830) [22], one started to distinguish between “function” and “analytic expression” for the function. The concept has evolved ever since. One can see a function as a subset \( Z \) of the Cartesian product \( X \times Y \) such that \( (\{x\} \times Y) \cap Z \) has exactly one point. A more algorithmic point of view emerged with the appearance of computer science pioneered by Fermat and Blaise Pascal (1623-1662). It is the concept of a function as a “rule” assigning to an “input” \( x \) an “output” \( y \). A notion of Turing machine made precise what it means to “compute” \( f(x) \). After the notion of analytic continuation has sunk in by Bernhard Riemann (1826-1866), it became clear that a function can be a world by itself with more information than anticipated. The zeta function \( \zeta(z) \) for example, given as a sum \( \zeta(s) = \sum_{n=1}^{\infty} n^{-s} \) can be extended beyond the domain \( Re(s) > 1 \), where the sum is defined to the entire complex plane except \( s = 1 \). We already know in school calculus that the geometric sum \( f(a) = \sum_{n=0}^{\infty} a^n \) makes sense for \( a > 1 \) if written as \( f(a) = 1/(1-a) \) defining so \( f(2) = -1 \) even if the sum \( \sum_{n} 2^n \) diverges. The theory of Riemann surfaces shows that functions often have an extended life as geometric objects. We have also learned that there are functions which are not computable. The function which assigns to a Turing machine the value 1 if it halts and 0 if not for example is not computable. The geometric point of view is how everything has started: functions were perceived as geometry very early on. The relation \( f(x, y, z) = x^2 + y^2 - z^2 = 0 \) for example would have been seen by Pythagoras as a relation of areas of three squares built from a right angle triangles and mathematicians like Brahmagupta (598-670) or Al-Khwarizmi (780-850) saw the construction of roots of the quadratic equation \( f(x) = x^2 + bx + c = 0 \) as a geometric problem about the area of some region which when suitably completed becomes a square.
### Appendix C: Mathematica code

Here are some of the basic functions:

| Function |
|----------|
| \( X[x_, n_] := \text{Product}[x - j, \{j, 0, n - 1\}] \) |
| \( \sin[a_, x_] := \text{Im}[\exp[a I, x]] \) |
| \( \cos[a_, x_] := \text{Re}[\exp[a I, x]] \) |
| \( \tan[a_, x_] := \sin[a, x] / \cos[a, x] \) |
| \( \cot[a_, x_] := \cos[a, x] / \sin[a, x] \) |
| \( \exp[a_, x_] := (1 + a h) \hat{(x/h)} \) |
| \( \log[a_, x_] := \text{Log}[1 + a x] \) |
| \( \text{rec}[x_] := \log[1 + x] - \log[1] \) |

Here is an illustration of the fundamental theorem of calculus:

| Code |
|------|
| \( d[f_] := \text{Function}[m, f[m + 1] - f[m]] \) |
| \( s[f_] := \text{Function}[m, \text{Sum}[f[k], \{k, 1, m\}]] \) |
| \( f0 = \text{Function}[x, \sin[1, x]] \) |
| \( g0 = d[f0] \) |
| \( h0 = s[g0] \) |
| \( \{f0[201], h0[201]\} \) |

And here is an example of a Taylor expansion leading to an interpolation:

| Code |
|------|
| \( f0[x_] := \cos[x \hat{2}/2] \) |
| \( d[f_] := \text{Function}[m, f[m + 1] - f[m]] \) |
| \( \text{Di}[f_, k_] := \text{Module}[\{g, h = f\}, \text{Do}[g = d[h]; h = g, \{j, k\}]; h] \) |
| \( X[x_, n_] := \text{Product}[x - j, \{j, 0, n - 1\}] \) |
| \( \text{Taylor}[f_] := \text{Table}[g = \text{Di}[f, k]; g[0], \{k, 0, M\}] \) |
| \( a = \text{Taylor}[f0] \) |
| \( f1[x_] := \text{Sum}[X[x, j] \ a[[j + 1]] / j!, \{j, 0, M\}] \) |
| \( S1 = \text{Plot}[\{f0[x], f1[x]\}, \{x, 0, M\}] \) |
| \( S2 = \text{Graphics}[\text{Table}[\{\text{Red, Point}[\{k, f1[k]\}], \{k, 0, M - 1\}\}]] \) |
| \( \text{Show}[\{S1, S2\}] \) |

The functions can also be defined using a more general “Planck constant” \( h \):

| Code |
|------|
| \( \exp[a_, h_, x_] := (1 + a h) \hat{(x/h)} \) |
| \( \cos[a_, h_, x_] := \text{Re}[\exp[I a h, x]] \) |
| \( \sin[a_, h_, x_] := \text{Im}[\exp[I a h, x]] \) |
| \( \text{Plot}[\{\cos[x], \cos[1, 0.1, x], \sin[x], \sin[1, 0.1, x]\}, \{x, 0, 4 Pi\}] \) |

Here are some graphs, already built into the computer algebra system:

| Code |
|------|
| \( \text{CompleteGraph}[5] \) |
| \( \text{CycleGraph}[7] \) |
| \( \text{WheelGraph}[8] \) |
| \( \text{StarGraph}[11] \) |
| \( \text{PD} = \text{PolyhedronData}; p = \text{PD["Platonic"]}; \) |
| \( \text{UG} = \text{UndirectedGraph} \) |
| \( \text{UG[Graph[PD[p[[3]], "SkeletonRules"]]} \) |
| \( \text{UG[Graph[PD[p[[4]], "SkeletonRules"]]} \) |
Figure 1. Basic functions $\sin, \cos, \tan, \cot, \exp, \log, x^2, x^3$ compared with the classical functions.

Figure 2. A comparison of $\sin_h(x), \sin(x)$ and $\cosh(x), \cos(x)$ for $h = 0.1$. 
IF ARCHIMEDES WOULD HAVE KNOWN FUNCTIONS ...

Figure 3. Interpolating the function \( \cos(x^2/2) \) using the discrete Taylor formula of Newton and Gregory.

Figure 4. Example graphs \( K_5, C_7, W_7, S_{10}, O, I \).

6. Calculus flavors

Calculus exists in many different flavors. We have infinite-dimensional versions of calculus like functional analysis and calculus of variations, calculus has been extended to less regular geometries and functions using the language of geometric measure theory, integral geometry or the theory of distributions. Calculus in the complex in the form of complex analysis has become its own field and calculus on more general spaces is known as differential topology, Riemannian geometry or algebraic geometry, depending on the nature of the objects and the choice of additional structures.
The term “quantum calculus” is used in many different contexts as mathematical objects can be “quantized” in different ways. One can discretize, replace commutative algebras with non-commutative algebras, or replace classical variational problems with path integrals. “Quantized calculus” is part of a more general quantum calculus, where a measure $\mu$ on the real line is chosen leading to the derivative $f' = \int (f(x+h) - f(x))/h \, d\mu(h)$. Quantized calculus is a special case, when $\mu$ is the Lebesgue measure and quantum calculus is the case when $\mu$ is a Dirac measure so that $f' = (f(x+h) - f(x))/h$. The measure $\mu$ defining the derivative has a smoothing effect if it is applied to space of derivatives. [12] defines quantized calculus as a derivative $df = [F,f]$, where $F : H \to H$ is a self-adjoint operator on a Hilbert space $H$ satisfying $F^2 = 1$. The Hilbert transform has this property. Real variables are replaced by self-adjoint operator, infinitesimals are compact operators $T$, the integral is the Dixmier trace $||T||$ and integrable functions are replaced by compact operators for which $||T|| = \lim(1/\log(n)) \sum_{k \leq n} \lambda_k$ exists in a suitable sense, where $\lambda_k$ are the eigenvalues of the operator $|T|$ arranged in decreasing order. The Dixmier trace neglects infinitesimals, compact operators $T$ for which $\lambda_k = o(k^{-1})$. A Fredholm module is a representation of an algebra $A$ as operators in the Hilbert space $H$ such that $[F,f]$ is infinitesimal (=compact) for all $f \in A$. In the simplest case, for single variable calculus, $F$ is the Hilbert transform.

It is evident that quantum calculus approaches like [12] (geared towards geometry) and [32] (geared towards number theory) would need considerable adaptations to be used in a standard calculus sequence for non-mathematicians. Quantum calculus flavors already exist in the form of business calculus, where the concept of limit is discarded. For a business person, calculus can be done on spreadsheets, integration is the process of adding up rows or columns and differentiation is looking at trends, the goal of course being to integrate up the changes and get trends for the future. Books like [30] (a “brief edition” with 800 pages) however follow the standard calculus methodology.

I personally believe that the standard calculus we teach today is and will remain an important benchmark theory which has a lot of advantages over discretized calculus flavors. From experience, I noticed for example that having to carry around an additional discretization parameter $h$ can be a turn-off for students. This is why in the talk, $h = 1$ was assumed. In the same way as classical mechanics is an elegant idealization of quantum mechanics and Newtonian mechanics is a convenient limit of general relativity, classical calculus is a teachable idealisation of a more general quantum calculus.

Still, a discrete approach to calculus could be an option for a course introducing proofs or research ideas. Many of the current calculus books have grown to huge volumes and resemble each other even to the point that different textbooks label chapters in the same order and that problems resemble each other, many of them appear generated by computer algebra systems.
But as many newcomers have seen, the transition from a standard textbook to a new textbook is hard. If there are too many changes even if it only affects the order in which the material is taught, then this can alienate the customer. Too fast innovation destroys its sales. Innovative approaches are seldom successful commercially, especially because college calculus systems have become complicated machineries, where larger courses with different sections are coordinated and teams of teachers and course assistants work together. When looking at the history of calculus textbooks from the past, over decades or centuries, we can see that the development today is still astonishingly fast today and that the variety has never been bigger.

A smaller inquiry based or proof based course aiming to have students in a “state of research” can benefit from fresh approaches. The reason is that asking to reprove standard material have become almost impossible today. It is due to the fact that one can look up proofs so easily today: just enter the theorem into a search engine and almost certainly a proof can be found. It is no accident that more inquiry based learning stiles have flourished in a time when point set topology has been developed. Early point set topology for example was a perfect toy to sharpen research skills. It has become almost impossible now to make any progress in basic set topology because the topic has been cataloged so well. When developing a new theorem, a contemporary student almost certainly will recover a result which is a special case of something which has been found 50-100 years ago. In calculus it is even more extreme because the theorems taught have been found 200-300 years ago.

Learning old material in a new form is always exciting. I myself have learned calculus the standard way but then seen it again in a “nonstandard analysis course” given by Hans Läuchli. The course had been based on the article [52]. We were also recommended to read [57], a book which stresses in the introduction how Euler has worked already in a discrete setting. Euler would write $e^x = (1 + x/n)^n$, where $n$ is a number of infinite size (Euler: “$n$ numerus infinite magnus”) ([21] paragraph 155). Nelson also showed that probability theory on finite sets works well ([53]. There are various textbooks which tried nonstandard analysis. One which is influenced by the Robinson approach to nonstandard analysis is [34]. But nonstandard analysis has never become the standard. It remains true to its name.

7. Calculus textbooks

The history of calculus textbooks is fascinating and would deserve to be written down in detail. Here are some comments about some textbooks from my own calculus library. The first textbook in calculus is L'Hôpital's “Analyse des Infiniment Petits pour l'Intelligence des Lignes Courbes” [14], a book from 1696, which had less than 200 pages and still contains 156 figures. Euler's textbook [21] on calculus starts the text with a modern treatment of functions but still emphasis "analytic expressions". Euler for example would hardly considered the rule which assigns 1 to an irrational numbers and 0 on rational number, to be a function. Euler uses already in that textbook already his formula $e^{ix} = \cos(x) + i \sin(x)$ extensively but in material is closer of what
we consider “precalculus” today, even so it goes pretty far, especially with series and continued fraction expansions. An other master piece is the textbook of Lagrange [15]. [33] mentions textbooks of the 18th century with authors C. Reyneau, G. Chayne, C. Hayes, H. Ditton, J. Craig, E. Stone, J. Hodgson, J. Muller or T. Simpson. Hardy’s book [28], whose first edition came out in 1908 also contains a substantial real analysis part and uses already the modern function approach. It no more restricts to “analytic expressions” but also allows functions to be limiting expressions of series of functions. It is otherwise a single variable text book. From about the same time is [50], which is more than 100 years old but resembles in style and size modern textbooks already, including lots of problems, also some of applied nature. The book [23] is a multivariable text which has many applications to physics and also contains some differential geometry. The book [46] is remarkable because it might be where we “stole” the first slide of our talk from, mentioning sticks and pebbles to make the first steps in calculus. This book weights only 150 pages but covers lots of ground. [55] treats single variable on 200 pages with a decent amount of exercises and [29] manages with the same size to cover a fair amount of multivariable calculus and tensor analysis, differential geometry and calculus of variations in three dimensions, uses illustrations however only sparingly. The structure of Hay’s book is in parts close to modern calculus textbooks like [56, 61, 63, 27], many of them started to grow fatter. The textbook [65] for example has already 420 pages but covers single and multivariable calculus, real analysis and measure theory, Fourier series and partial differential equations. But it is no match in size to books like [27] with 1318 pages or Stewart with about the same size. Important in Europe was [13] which had been translated into other language and which my own teachers still recommended. Its influence is clearly visible in all major textbooks which are in use today. [47] made a big impact as it is written with great clarity.

Since we looked at a discrete setting, we should mention that calculus books based on differences have appeared already 60 years ago. The book [31] illustrates the difficulty to find good notation. This is a common obstacle with discrete approaches, especially also with numerical analysis books. Discrete calculus often suffers from ugly notation using lots of indices. One of the reason why Leibnitz calculus has prevailed is the use of intuitive notation. For a modern approach see [26].

Not all books exploded in size. [24] covers in 120 pages quite a bit of real analysis and multivariable calculus. But books would become heavier. [2] already counted 725 pages. It was intended for a three semester course and stated in the introduction: there is an element of truth in the old saying that the Euler textbook ‘Introductio in Analysis Infinitorum, (Lausannae, 1748)’ was the first great calculus textbook and that all elementary calculus textbooks copied from books that were copied from Euler. The classic [35] had already close to 1000 pages and already included a huge collection of problems. Influential for the textbooks of our modern time are also Apostol [3, 4] which I had used myself to teach. It uses elements of linear algebra, probability theory, real analysis, complex numbers and differential equations. Its influence is also clearly imprinted on other textbooks. The book [19] was one of the first which dared to
introduce calculus in arbitrary dimensions using differential forms. Since then, calculus books have diversified and exploded in general in size. There are also exceptions: original approaches include [59, 60, 51, 48, 54, 51, 44, 20, 25, 36].

8. Tweets

We live in an impatient twitter time. Here are two actually tweeted 140 character statements which together cover the core of the fundamental theorem of calculus:

| With $s(x) = x + 1$
| and inverse $t(x) = x - 1$,
| define $X^n = (xt)^n$.
| For example,
| $1 = 1$, $X = x$,
| $X^2 = x(x - 1)$,
| $X^3 = x(x - 1)(x - 2)$.

and

| $Df(x) = [s, f] = f(x + 1) - f(x)$
| gives $DX^n = nX^{n-1}$. With
| $Sf(x) = f(0) + f(1) + \ldots + f(n - 1)$,
| the fundamental theorem
| $SDf(x) = f(x) - f(0)$
| $DSf(x) = f(x)$
| holds.

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