Detection of Vampire Attacks in Ad-Hoc Wireless Sensor Network Evaluation and Protection

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ABSTRACT
Wireless Sensor Networks (WSNs) in today’s world are the means of communication. These contain nodes that act as transmitter and receivers are prone to different attacks leading to different types of losses. The resource depletion attack that is called vampire attack drains out the energy from the nodes leaving them useless. These attacks are protocol compliant, they are easy to implement. Since they are orthogonal in nature they can easily intrude into any routing protocol. They affect the entire network causing large loss of energy and A vampire attack is caused by the malicious node on the decentralized ad hoc wireless network. The paper analyses how protocols faces these attacks. Vampire attacks are not protocol specific rather uses its compliant message. The current security measures to prevent these attacks are been reviewed along with result of simulation of representative protocols in the presence of a vampire attack is been presented. The paper also describes how the existing sensor network protocol is been modified for protection from the vampire attacks for which PLGP (Parno, Luk, Gaustad and Perrig) solution is also been proposed.

KEYWORDS: Wireless Sensor Network; Denial of service; Resource depletion; Routing; Energy consumption; Security; carousel attack; stretch attack; PLGP

I. INTRODUCTION
Ad hoc wireless sensor network consists of various sensors that are expanded in a space where each sensor performs signal processing and data networking providing operational efficiency. The ad hoc wireless servers are self-organized and energy constrained. These sensor networks are used to detect information of enemy base, monitor environmental changes and are also used for security purposes in various places like shopping and parking area and when these networks face attacks causing negative effect by causing battery exhaustion and higher energy utilization.

Fig 1: Ad hoc Wireless sensor network
Vampire attacks are caused when a message is been initiated and transmitted through a malicious node over the network causing higher battery utilization and battery exhaustion. Vampire attacks are not constrained to a specific type of protocol and does not alter specific path in the network. When a network is been attacked by them, even transfer of small data consumes more energy.

Fig 2: Architecture of WSN
The basic architecture of a WSN [10] contains number of nodes that act as routing nodes, sensing nodes and base nodes. Figure 2 shows the architecture of a WSN, where all the different types of nodes are shown. It is also possible that every node in the network functions as an aggregation point, delaying transmission of an outgoing message until a sufficient number of incoming messages have been received and aggregated. Power management in sensor networks is critical. Consequently, if we want sensor networks to last for years, it is crucial that they run at around a 1% duty cycle (or less). Similarly, since the power consumption of the radio is three orders of magnitude higher when transmitting or listening than when in sleep mode, it is crucial to keep the radio in sleep mode the overwhelming majority of the time. Sensor networks provide economically viable solutions for a wide variety of applications, including surveillance of critical infrastructure, safety monitoring, and many health-care applications [6]. As sensor networks are increasingly deployed in such security-and safety critical environments, the need for secure communication primitives is self-evident. Likewise, the development of such primitives facilitates the use of sensor networks in a wider range of applications. The central goal of this work is to ensure node-to-node message delivery, even if the sensor network is under active attack.
great as the ones contained in any RREQ that the node has received for that destination) it broadcasts a route request (RREQ) message to its neighbors. Each node that receives the broadcast sets up a reverse route towards the originator of the RREQ (unless it has a “fresher” one). When the intended destination (or an intermediate node that has a “fresh enough” route to the destination) receives the RREQ, it replies by sending a Route Reply (RREP). It is important to note that the only mutable information in a RREQ and in a RREP is the hop count (which is being monotonically increased at each hop). The RREP travels back to the originator of the RREQ (this time as a unicast). At each intermediate node, a route to the destination is set (again, unless the node has a “fresher” route than the one specified in the RREP). In the case that the RREQ is replied to by an intermediate node (and if the RREQ had set this option), the intermediate node also sends a RREP to the destination. In this way, it can be granted that the route path is being set up bidirectional. In the case that a node receives a new route (by a RREQ or by a RREP) and the node already has a route “as fresh” as the received one, the shortest one will be updated. If there is a subnet (a collection of nodes that are identified by a common network prefix) that does not use AODV as its routing protocol and wants to be able to exchange information with an AODV network, one of the nodes of the subnet can be selected as their “network leader”. The network leader is the only node of the subnet that sends forwards and processes AODV routing messages and every RREP that the leader issues, it sets the prefix size of the subnet optionally. A Route Reply Acknowledgment (RREP-ACK) message may be sent by the originator of the RREQ to acknowledge the receipt of the RREP.
III. VAMPIRE ATTACK DETECTION

There are two types of attacks in WSN, the routing depletion and resource depletion attack. The routing depletion attacks usually only affect the routing path the resource depletion attacks are the ones that attack the network features like bandwidth, power, and energy consumption. These attacks are commonly called as “Vampire attacks” [2]. They are called so because they drain the battery power from the nodes. These are a type of Denial of Service [1] since they affect the entire system from performing. They are difficult to be detected since they are protocol compliant and are orthogonal to them [4]. They are not protocol specific. They do not affect a single node they take their time attack one by one and disrupt the entire system. Vampire attacks can be defined as the composition and transmission of a message that cause more energy to be consumed by the network than if an honest node transmitted a message of identical size to the same destination, although using different packet headers. The strength of the attack is measured by the ratio of network energy used in the benign case to the energy used in the malicious case. Safety from Vampire attacks implies that this ratio is 1. Energy use by malicious nodes is not considered, since they can always unilaterally drain their own batteries.

A. Carousel attack

In this attack, an adversary composes packets with purposely introduced routing loops. It is called carousel attack, since it sends packets in circles as shown in Figure 4. It targets source routing protocols by exploiting the limited verification of message headers at forwarding nodes, allowing a single packet to repeatedly traverse the same set of nodes. On average, a random located carousel attacker in the example mentioned topology can increase the network energy consumption by a factor of 1.48±0.99.

B. Stretch attack

In this attack, also targeting source routing, an adversary constructs artificially long routes, potentially traversing every node in the network. It is call this the stretch attack, since it increases packet path lengths, causing packets to be processed by a number of nodes that is independent of hop count along the shortest path between the adversary and packet destination. An example is illustrated in figure 6. In the example topology, there is an increase in energy usage by as much as a factor of 10.5 per message over the honest scenario, with an average increases in energy consumption of 2.67±2.49. As with the carousel attack, the reason for the large standard deviation is that the position of the adversarial node affects the strength of the attack. Not all routes can be significantly lengthened, depending on the location of the adversary.

IV. SECURITY AGAINST VAMPIRE ATTACKS

A Clean Slate Sensor Network Routing by PLGP (Parno, Luk, Gaustad and Perrig) can be applied which consists of two phases:

1. Topology Discovery Phase
2. Packet Forwarding Phase

1. Topology Discovery Phase:

A node starts with its virtual address as zero. At each node a certificate is been issues which contains the public key for identification. Each node is connected to the other and shares virtual address, public key and the certificate when they merge with closest nearby group.
2. Packet Forwarding Phase:
The packets are forwarded in this phase as shown in figure 8.

PLGP proposed a solution which suggests:
A. Providing a verifiable path history to all the packets involved.
B. Using this path history the packet transmission can take place through every node securely passing through at least one honest node.
C. Each node upon receiving the message, checks for authentication in the chain.

CONCLUSION
In this paper, the routing protocol affected by vampire attack in WSN is discussed. This is a new class of resource consumption attack that use routing protocols to permanently disable ad-hoc WSNs by depleting node’s battery power. Simulation results show that depending on the location of adversary, network energy expenditure during the forwarding phase increasing. The security flaws of AODV can be fixed by using RSA encryption system that will avoid the adversary from entering the system. These attacks does not depend on particular type of protocol and Ad hoc network sensors have been applied in various fields which needs to create and identify solutions for prevention of the network from these attacks. There are different types of vampire attacks depending on the protocol. When the attack take place it not only consumes higher power but also takes additional time. There are many solutions and techniques that have been presented to prevent these attacks but were not effective enough which creates a need for a better solution. PLGP solutions can be applied to these protocols in order to prevent these networks that are often prone to vampire attacks.

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