ON THE ORIENTED CHROMATIC NUMBER
OF DENSE GRAPHS

DAVID R. WOOD

Abstract. Let $G$ be a graph with $n$ vertices, $m$ edges, average degree $\delta$, and maximum degree $\Delta$. The oriented chromatic number of $G$ is the maximum, taken over all orientations of $G$, of the minimum number of colours in a proper vertex colouring such that between every pair of colour classes all edges have the same orientation. We investigate the oriented chromatic number of graphs, such as the hypercube, for which $\delta \geq \log n$. We prove that every such graph has oriented chromatic number at least $\Omega(\sqrt{n})$. In the case that $\delta \geq (2+\epsilon) \log n$, this lower bound is improved to $\Omega(\sqrt{m})$. Through a simple connection with harmonious colourings, we prove a general upper bound of $O(\Delta \sqrt{n})$ on the oriented chromatic number. Moreover this bound is best possible for certain graphs. These lower and upper bounds are particularly close when $G$ is $(c \log n)$-regular for some constant $c > 2$, in which case the oriented chromatic number is between $\Omega(\sqrt{n} \log n)$ and $O(\sqrt{n} \log n)$.

1. Introduction

Throughout this paper, $G$ is a (finite and simple) undirected graph with $n$ vertices, $m$ edges, and maximum degree $\Delta$. A colouring of $G$ is a function that assigns a ‘colour’ to each vertex so that adjacent vertices receive distinct colours. The chromatic number $\chi(G)$ is the minimum number of colours in a colouring of $G$. An oriented graph $D$ is a directed graph with no parallel and no antiparallel arcs; that is, no two arcs have the same pair of endpoints. If $G$ is the underlying undirected graph of $D$ then $D$ is an orientation of $G$. An oriented colouring of $D$ is a colouring of $G$ such that between each pair of colour classes, all edges have the same direction. That is, there are no arcs $\vec{v}w$ and $\vec{xy}$ with $c(v) = c(y)$ and $c(w) = c(x)$. The oriented chromatic number $\overrightarrow{\chi}(D)$ is the minimum number of colours in an oriented colouring of $D$. The oriented chromatic number $\overrightarrow{\chi}(G)$ is the maximum of $\overrightarrow{\chi}(D)$, taken over all orientations $D$ of $G$. The oriented chromatic number was introduced by Courcelle in 1994 and is now a widely studied parameter; see \cite{3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 21, 22, 23, 24, 25, 26, 27, 28, 29, 30, 31, 32, 33, 34, 35, 36, 37, 38, 39, 40, 41, 42, 43, 44, 45, 46, 47, 48, 49, 50, 51, 52, 53, 54, 55, 56, 57}.

This paper is motivated by a question of André Raspaud [private communication, Prague 2004], who asked for the oriented chromatic number of the $d$-dimensional hypercube $Q_d$. This

\textbf{Date:} August 2, 2021.

\textbf{2000 Mathematics Subject Classification.} 05C15 (coloring of graphs and hypergraphs).

\textbf{Key words and phrases.} graph, graph colouring, oriented colouring, oriented chromatic number, hypercube, harmonious colouring.

Supported by a Marie Curie Fellowship of the European Community under contract 023865, and by the projects MCYT-FEDER BFM2003-00368 and Gen. Cat 2001SGR00224.

1
is the graph with vertex set \( \{0,1\}^d \), where two vertices are adjacent whenever they differ in precisely one coordinate. \( Q_d \) is \( d \)-regular and has \( 2^d \) vertices. In this paper we prove generally applicable bounds on \( \overrightarrow{\chi}(G) \), which in the case of the hypercube give

\[
0.8007... \sqrt{2^d} \leq \overrightarrow{\chi}(Q_d) \leq 2d\sqrt{2^d - 1},
\]

thus determining \( \overrightarrow{\chi}(Q_d) \) to within a factor of about \( \frac{5}{2} \). No non-trivial bounds on \( \overrightarrow{\chi}(Q_d) \) were previously known, as we now describe.

An undirected graph has \( \chi(G) = n \) if and only if \( G = K_n \); that is, if \( G \) has diameter 1. But when does \( \overrightarrow{\chi}(G) = n \)? This question was asked by Füredi et al. [25], who observed that for every oriented graph \( D \),

\[
\overrightarrow{\chi}(G) = n \text{ if and only if } D \text{ has diameter 2.}
\]

Here the \textit{diameter} of \( D \) is the least integer \( k \) such that every pair of vertices in \( D \) are connected by a directed path of at most \( k \) edges. Klostermeyer and MacGillivray [33] called an oriented graph with diameter 2 an \textit{oclique}. Note that small diameter (\( > 2 \)) does not necessarily imply large oriented chromatic number. For example, \( K_{1,1,n} \) has an orientation with diameter 3 and oriented chromatic number 3. Erdős et al. [20] proposed studying the extremal function \( f(n) \), defined to be the minimum number of arcs in an oriented graph with \( n \) vertices and diameter 2. Katona and Szemerédi [32] proved that \( \frac{n}{2} \log \frac{n}{2} \leq f(n) \leq n \lceil \log n \rceil \), and Füredi et al. [25] tightened both bounds to conclude that \( f(n) = (1 - o(1))n \log n \). These results imply that there are \( n \)-vertex graphs with the same number of edges as the hypercube (that is, \( n \log n \)), yet have oriented chromatic number \( n \). Thus good bounds for \( \overrightarrow{\chi}(Q_d) \) cannot be obtained just in terms of the number of edges.

The example of an oriented graph with diameter 2 by Füredi et al. [25] has a vertex of degree \( n - 1 \). Thus it is natural to consider the oriented chromatic number of graphs with bounded degree. Sopena [52] and Kostochka et al. [37] proved that the oriented chromatic number is bounded for graphs of bounded degree. The best bound is due to Kostochka et al. [37], who proved that every graph \( G \) satisfies

\[
\overrightarrow{\chi}(G) \leq 2\Delta 2^{\Delta},
\]

and if \( G \) is \( \Delta \)-regular with sufficiently many vertices then

\[
\overrightarrow{\chi}(G) \geq 2^{\Delta/2}.
\]

Thus the exponential dependence on \( \Delta \) in (3) is unavoidable. Observe that for graphs such as the hypercube with \( \Delta \geq \log n \), the upper bound in (3) is greater than the trivial upper bound of \( n \).

This motivates the study of the oriented chromatic number of graphs whose average degree is at least logarithmic in the number of vertices. For any such graph we establish a lower bound of \( \Omega(\sqrt{n}) \) on the oriented chromatic number. If the average degree is at least \( (2 + \epsilon) \log n \) then this lower bound is improved to \( \Omega(\sqrt{n}) \). These results are proved in Section 3. In Section 2 we use a simple connection with harmonious colourings to prove a general upper bound of
$O(\Delta \sqrt{n})$ on the oriented chromatic number. Moreover this bound is best possible for certain graphs, as proved in Section 3.

2. An Upper Bound

In this section we prove an elementary upper bound on the oriented chromatic number. A colouring of an undirected graph $G$ is harmonious if the endpoints of every pair of distinct edges receive at least three colours. That is, every bichromatic subgraph has at most one edge. The harmonious chromatic number $h(G)$ is the minimum number of colours in a harmonious colouring of $G$; see [2, 8, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 26, 27, 28, 29, 38, 39, 39, 40, 42, 43, 56].

The following two basic lower bounds on $h(G)$ are well known. Since a vertex and its neighbours all receive distinct colours in a harmonious colouring,

$$h(G) \geq \Delta + 1.$$

Since $G$ has at most $\left(\frac{h(G)}{2}\right)$ edges,

$$(5) \quad h(G) > \sqrt{2m}.$$

The next bound is new. Observe that a harmonious colouring of $G$ is an oriented colouring for every orientation of $G$. Thus

$$(6) \quad \overrightarrow{\chi}(G) \leq h(G).$$

Many upper bounds on $h(G)$ are known. For example, McDiarmid and Xinhua [42] proved that $h(G) \leq 2\Delta \sqrt{n - 1}$. Thus (6) implies the following lemma, which proves the upper bound on $\overrightarrow{\chi}(Q_d)$ in (1).

**Lemma 1.** For every graph $G$,

$$\overrightarrow{\chi}(G) \leq 2\Delta \sqrt{n - 1}.$$

3. An Existential Lower Bound

In this section we construct a graph whose oriented chromatic number is within a constant factor of the upper bound in Lemma 1.

**Lemma 2.** For infinitely many $n$, there is an $n$-vertex $\Delta$-regular graph $G$ with $\overrightarrow{\chi}(G) = n$ and $\Delta = \sqrt{8n + 1} - 3$.

**Proof.** As illustrated in Figure 1, let $G$ be the underlying undirected graph of the oriented graph $D$ with vertex set $\{(i, j) : 1 \leq i < j \leq p\}$ and arcs

(a) $(i, j)(i, k)$ whenever $i < j < k$,
(b) $(i, j)(k, j)$ whenever $i < k < j$, and
(c) $(i, j)(k, i)$ whenever $k < i < j$.

First we compute the degree of each vertex $(i, j)$. Observe that $(i, j)$ has $p - j$ outgoing type-(a) arcs, $j - i - 1$ outgoing type-(b) arcs, and $i - 1$ outgoing type-(c) arcs. Thus $(i, j)$
has outdegree \((p - j) + (j - i - 1) + (i - 1) = p - 2\). A type-(a) incoming arc at \((i, j)\) is from a vertex \((i, k)\) with \(i < k < j\); there are \(j - i - 1\) such arcs. A type-(b) incoming arc at \((i, j)\) is from a vertex \((k, j)\) with \(k < i < j\); there are \(i - 1\) such arcs. A type-(c) incoming arc at \((i, j)\) is from a vertex \((j, k)\) with \(i < j < k\); there are \(p - j\) such arcs. Thus \((i, j)\) has outdegree \((j - i + 1) + (i - 1) + (p - j) = p - 2\). Hence \(G\) is \(\Delta\)-regular with \(\Delta = 2(p - 2) = \sqrt{8n + 1} - 3\).

Suppose on the contrary that \(D\) has a directed 2-cycle \(C\). If \(C\) has a type-(a) arc \((i, j)(i, k)\), then the reverse arc \((i, k)(i, j)\) is also type-(a), implying \(j < k\) and \(k < j\), which is a contradiction. If \(C\) has a type-(b) arc \((i, j)(k, j)\), then the reverse arc \((k, j)(i, j)\) is also type-(b), implying \(k < j\) and \(j < k\), which is a contradiction. If \(C\) has a type-(c) arc \((i, j)(k, i)\), then the reverse arc is \((k, i)(i, j)\), but there are no arcs of this form. Thus \(D\) has no directed 2-cycle, and indeed \(D\) is an oriented graph.

We claim that \(D\) has diameter 2. Consider two vertices \((i, j)\) and \((k, \ell)\). Then \(i < j\) and \(k < \ell\), and without loss of generality, \(i \leq k\). If \(i = k\) and \(j < \ell\), then \((i, j)(i, \ell)\) is a type-(a) arc of \(D\). If \(i = k\) and \(\ell < j\), then \((i, \ell)(i, j)\) is a type-(a) arc of \(D\). Now assume that \(i < k\). If \(i < k\) and \(j = \ell\), then \((i, j)(k, j)\) is a type-(b) arc of \(D\). If \(i < k\) and \(j < \ell\), then \((i, j)(i, \ell)(k, \ell)\) is a type-(ab) path of \(D\). Otherwise \(i < k\) and \(\ell < j\), implying \(i < k < \ell < j\), in which case \((i, j)(\ell, j)(k, \ell)\) is a type-(bc) path of \(D\). Thus \(D\) has diameter 2, implying \(\overrightarrow{\chi}(D) = n\) by [2].

It follows from Lemma [2] that in any upper bound of the form \(\overrightarrow{\chi}(G) \leq O(\Delta^\alpha n^\beta)\), we must have \(\alpha + 2\beta \geq 2\). In particular with \(\beta = \frac{1}{2}\), for the graph \(G\) from Lemma [2] we have \(\overrightarrow{\chi}(G) = n > \Delta \sqrt{\frac{n}{8}}\). In this sense, the upper bound in Lemma [1] is tight up to a constant factor.

Also note that Moore’s bound for the degree/diameter problem implies that \(\Delta \geq \sqrt{n - 1}\) in every (undirected) graph with diameter 2.
4. A Universal Lower Bound

We now consider universal lower bounds on the oriented chromatic number. Kostochka et al. [37] proved the following lower bound for all $G$, which implies (4). (Throughout this paper, all logarithms are binary.)

$$\left(\frac{\chi^\rightarrow(G)}{2}\right) + n \log(\chi^\rightarrow(G)) \geq m.$$  

We now reformulate (7) for reasonably dense graphs. Say $G$ has average degree $\delta := \frac{2m}{n}$. Let $t$ be the solution to

$$t + \log t = \delta - \log n.$$  

Note that $0 < t < \delta$ and $t \to \delta$ for $\delta \gg \log n$. (We are not interested in the case $\delta \ll \log n$, when $t$ becomes small.)

**Lemma 3.** For every graph $G$,

$$\chi^\rightarrow(G) \geq \sqrt{nt}.$$  

**Proof.** Suppose on the contrary that $\chi^\rightarrow(G) < \sqrt{nt}$. By (7),

$$\left(\frac{\sqrt{nt}}{2}\right) + n \log(\sqrt{nt}) > m.$$  

Thus $\frac{1}{2}nt + \frac{1}{2}n \log(nt) > \frac{1}{2}\delta n$, implying $t + \log t > \delta - \log n$. This contradiction proves the claim. \qed

**Lemma 4.** For every graph $G$ with average degree $\delta \geq \log n$,

$$\chi^\rightarrow(G) \geq 0.8007... \sqrt{n}.$$  

**Proof.** Lemma 3 implies the claim since $\sqrt{t} \geq 0.8007...$ whenever $\delta \geq \log n$. \qed

For the hypercube, $\delta = \log n$. Thus Lemma 3 implies the lower bound in (11). Since (7) is proved by a non-constructive counting argument, it would be interesting to construct an orientation $D$ of $Q_d$ with $\chi^\rightarrow(D) \in \Omega(\sqrt{2^d})$; see [1, 22, 24, 41] for results on specific orientations of the hypercube.

We now refine Lemma 3 for graphs that are more dense than hypercubes.

**Lemma 5.** For every graph $G$ with average degree $\delta \geq \log n + (1 + \epsilon) \log t$ for some $\epsilon > 0$,

$$\chi^\rightarrow(G) \geq \sqrt{\frac{\epsilon}{1+\epsilon} \left(2m - n \log n\right)}.$$  

(For example, the assumption in Lemma 5 holds if $\delta \geq (2 + \epsilon) \log n$.)

---

4For completeness we include the proof of Equation (7) by Kostochka et al. [37]. Let $k := \chi^\rightarrow(G)$. $G$ has less than $k^n$ colourings with $k$ colours, each of which is an oriented colouring of at most $2 \binom{k}{2}$ orientations. Thus the number of orientations, $2^m$, is less than $k^n 2 \binom{k}{2}$. Thus $m < n \log k + \binom{k}{2}$.
Proof. By the assumption, \( t + \log t \geq (1 + \epsilon) \log t \) and \( t > \epsilon \log t \). Thus
\[
(1 + \epsilon)t > \epsilon(t + \log t) = \epsilon(\delta - \log n) \quad \text{and} \quad (1 + \epsilon)tn > \epsilon(\delta n - n \log n) = \epsilon(2m - n \log n).
\]
Therefore Lemma 3 implies that
\[
\chi'(G) \geq \sqrt{tn} > \sqrt{\frac{1}{1+\epsilon}} (2m - n \log n).
\]

Lemma 5 says that for sufficiently dense graphs (that is, graphs with super-logarithmic average degree) the lower bound of \( h(G) \geq \Omega(\sqrt{m}) \) in [4] also holds for the oriented chromatic number.

Now suppose that \( G \) is \( \Delta \)-regular for some \( \Delta \geq (2 + \epsilon) \log n \). Thus Lemmas 1 and 5 determine \( \chi'(G) \) to within a factor of \( \Theta(\sqrt{\Delta}) \). In particular,
\[
\sqrt{\frac{1}{2+\epsilon}} \Delta n \leq \sqrt{\frac{1}{2+\epsilon}} (\Delta - \log n)n \leq \chi'(G) \leq 2\Delta \sqrt{n-1}.
\]
The bounds in (8) are particularly close when \( G \) is \((c \log n)\)-regular for some constant \( c > 2 \). Then \( \chi'(G) \) is between \( \Omega(\sqrt{n \log n}) \) and \( O(\sqrt{n \log n}) \).

Acknowledgements

Thanks to Vida Dujmovi´c, Andr´e Raspaud, Bruce Reed, J´eан-Sebastien Sereni, Ricardo Strausz, and St´ephan Thomassé for stimulating discussions.

References

[1] Pierre Baldi. Neural networks, orientations of the hypercube, and algebraic threshold functions. IEEE Trans. Inform. Theory, 34(3):523–530, 1988.
[2] Donald G. Beane, Norman L. Biggs, and Brian J. Wilson. The growth rate of the harmonious chromatic number. J. Graph Theory, 13(3):291–299, 1989.
[3] Oleg V. Borodin, Dmitry Fon-Der-Flaass, Alexandr V. Kostochka, Andr´e Raspaud, and Éric Sopena. On deeply critical oriented graphs. J. Combin. Theory Ser. B, 81(1):150–155, 2001.
[4] Oleg V. Borodin and A. O. Ivanova. Oriented 7-coloring of planar graphs with girth at least seven. Sib. `Elektron. Mat. Izv., 2:222–229, 2005.
[5] Oleg V. Borodin and A. O. Ivanova. Oriented coloring of planar graphs with girth at least four. Sib. `Elektron. Mat. Izv., 2:239–249, 2005.
[6] Oleg V. Borodin, Alexandr V. Kostochka, Jaroslav Neˇsetˇril, Andr´e Raspaud, and Éric Sopena. On universal graphs for planar oriented graphs of a given girth. Discrete Math., 188(1-3):73–85, 1998.
[7] Oleg V. Borodin, Alexandr V. Kostochka, Jaroslav Neˇsetˇril, Andr´e Raspaud, and Éric Sopena. On the maximum average degree and the oriented chromatic number of a graph. Discrete Math., 206(1-3):77–89, 1999.
[8] Duncan Campbell and Keith Edwards. A new lower bound for the harmonious chromatic number. *Australas. J. Combin.*, 29:99–102, 2004.

[9] Bruno Courcelle. The monadic second order logic of graphs. VI. On several representations of graphs by relational structures. *Discrete Appl. Math.*, 54(2-3):117–149, 1994.

[10] Keith Edwards. The complexity of some graph colouring problems. *Discrete Appl. Math.*, 36(2):131–140, 1992.

[11] Keith Edwards. The harmonious chromatic number of almost all trees. *Combin. Probab. Comput.*, 4(1):31–46, 1995.

[12] Keith Edwards. The harmonious chromatic number of bounded degree trees. *Combin. Probab. Comput.*, 5(1):15–28, 1996.

[13] Keith Edwards. The harmonious chromatic number and the achronatic number. In *Surveys in combinatorics*, vol. 241 of *London Math. Soc. Lecture Note Ser.*, pp. 13–47. Cambridge Univ. Press, 1997.

[14] Keith Edwards. The harmonious chromatic number of bounded degree graphs. *J. London Math. Soc. (2)*, 55(3):435–447, 1997.

[15] Keith Edwards. A new upper bound for the harmonious chromatic number. *J. Graph Theory*, 29(4):257–261, 1998.

[16] Keith Edwards. The harmonious chromatic number of complete r-ary trees. *Discrete Math.*, 203(1-3):83–99, 1999.

[17] Keith Edwards. Detachments of complete graphs. *Combin. Probab. Comput.*, 14(3):275–310, 2005.

[18] Keith Edwards and Colin McDiarmid. New upper bounds on harmonious colorings. *J. Graph Theory*, 18(3):257–267, 1994.

[19] Keith Edwards and Colin McDiarmid. The complexity of harmonious colouring for trees. *Discrete Appl. Math.*, 57(2-3):133–144, 1995.

[20] Paul Erdős, Alfred Rényi, and Vera T. Sós. On a problem of graph theory. *Studia Sci. Math. Hungar.*, 1:215–235, 1966.

[21] Louis Esperet and Pascal Ochem. Oriented colorings of 2-outerplanar graphs. *Information Processing Letters*, to appear.

[22] Hazel Everett and Arvind Gupta. Acyclic directed hypercubes may have exponential diameter. *Inform. Process. Lett.*, 32(5):243–245, 1989.

[23] Guillaume Fertin, André Raspaud, and Arup Roychowdhury. On the oriented chromatic number of grids. *Inform. Process. Lett.*, 85(5):261–266, 2003.

[24] Pierre Fraigniaud, Jean-Claude König, and Emmanuel Lazard. Oriented hypercubes. *Networks*, 39(2):98–106, 2002.

[25] Zoltán Füredi, Peter Horak, Chandra M. Pareek, and Xuding Zhu. Minimal oriented graphs of diameter 2. *Graphs Combin.*, 14(4):345–350, 1998.

[26] John P. Georges. On the harmonious coloring of collections of graphs. *J. Graph Theory*, 20(2):241–254, 1995.
[27] Johannes H. Hattingh, Michael A. Henning, and Elna Ungerer. Graphs with small upper line-distinguishing and upper harmonious chromatic numbers. *J. Combin. Math. Combin. Comput.*, 47:165–181, 2003.

[28] Johannes H. Hattingh, Michael A. Henning, and Elna Ungerer. Upper line-distinguishing and upper harmonious chromatic numbers of cycles. *J. Combin. Math. Combin. Comput.*, 45:137–144, 2003.

[29] John E. Hopcroft and Mukkai S. Krishnamoorthy. On the harmonious coloring of graphs. *SIAM J. Algebraic Discrete Methods*, 4(3):306–311, 1983.

[30] Mohammad Hosseini Dolama and Eric Sopena. On the oriented chromatic number of graphs with given excess. *Discrete Math.*, 306(13):1342–1350, 2006.

[31] Mohammad Hosseini Dolama and Eric Sopena. On the oriented chromatic number of Halin graphs. *Inform. Process. Lett.*, 98(6):247–252, 2006.

[32] Gyula O. H. Katona and Endre Szemerédi. On a problem of graph theory. *Studia Sci. Math. Hungar*, 2:23–28, 1967.

[33] William F. Klostermeyer and Gary MacGillivray. Analogues of cliques for oriented coloring. *Discuss. Math. Graph Theory*, 24(3):373–387, 2004.

[34] William F. Klostermeyer and Gary MacGillivray. Homomorphisms and oriented colorings of equivalence classes of oriented graphs. *Discrete Math.*, 274(1-3):161–172, 2004.

[35] William F. Klostermeyer and Gary MacGillivray. Pushing vertices and oriented colorings. *Bull. Inst. Combin. Appl.*, 40:49–58, 2004.

[36] Alexandr V. Kostochka, Tomasz Luczak, Gábor Simonyi, and Eric Sopena. On the minimum number of edges giving maximum oriented chromatic number. In *Contemporary trends in discrete mathematics*, vol. 49 of *DIMACS Ser. Discrete Math. Theoret. Comput. Sci.*, pp. 179–182. Amer. Math. Soc., 1999.

[37] Alexandr V. Kostochka, Eric Sopena, and Xuding Zhu. Acyclic and oriented chromatic numbers of graphs. *J. Graph Theory*, 24(4):331–340, 1997.

[38] Ilia Krasikov and Yehuda Roditty. Bounds for the harmonious chromatic number of a graph. *J. Graph Theory*, 18(2):205–209, 1994.

[39] Marek Kubale. Harmonious coloring of graphs. In *Graph colorings*, vol. 352 of *Contemp. Math.*, pp. 95–104. Amer. Math. Soc., 2004.

[40] Sin-Min Lee and John Mitchem. An upper bound for the harmonious chromatic number of a graph. *J. Graph Theory*, 11(4):565–567, 1987.

[41] Jiří Matoušek. The number of unique-sink orientations of the hypercube. *Combinatorica*, 26(1):91–99, 2006.

[42] Colin McDiarmid and Luo Xinhua. Upper bounds for harmonious colorings. *J. Graph Theory*, 15(6):629–636, 1991.

[43] Zevi Miller and Dan Pritikin. The harmonious coloring number of a graph. *Discrete Math.*, 93(2-3):211–228, 1991.
[44] Jaroslav Nešetřil and André Raspaud. Antisymmetric flows and strong colourings of oriented graphs. *Ann. Inst. Fourier (Grenoble)*, 49(3):1037–1056, 1999.

[45] Jaroslav Nešetřil, André Raspaud, and Éric Sopena. Colorings and girth of oriented planar graphs. *Discrete Math.*, 165/166:519–530, 1997.

[46] Jaroslav Nešetřil, Éric Sopena, and Laurence Vignal. $T$-preserving homomorphisms of oriented graphs. *Comment. Math. Univ. Carolin.*, 38(1):125–136, 1997.

[47] Pascal Ochem. Oriented colorings of triangle-free planar graphs. *Inform. Process. Lett.*, 92(2):71–76, 2004.

[48] André Raspaud and Éric Sopena. Good and semi-strong colorings of oriented planar graphs. *Inform. Process. Lett.*, 51(4):171–174, 1994.

[49] Attila Sali and Gábor Simonyi. Oriented list colouring of undirected graphs. In *Contemporary trends in discrete mathematics*, vol. 49 of *DIMACS Ser. Discrete Math. Theoret. Comput. Sci.*, pp. 307–316. Amer. Math. Soc., 1999.

[50] Robert Šámal. Antisymmetric flows and strong oriented coloring of planar graphs. *Discrete Math.*, 273(1-3):203–209, 2003.

[51] Robert Šámal. Antiflows, oriented and strong oriented colorings of graphs. *Arch. Math. (Brno)*, 40(4):335–343, 2004.

[52] Éric Sopena. The chromatic number of oriented graphs. *J. Graph Theory*, 25(3):191–205, 1997.

[53] Éric Sopena. Oriented graph coloring. *Discrete Math.*, 229(1-3):359–369, 2001.

[54] Éric Sopena. There exist oriented planar graphs with oriented chromatic number at least sixteen. *Inform. Process. Lett.*, 81(6):309–312, 2002.

[55] Andrzej Szepietowski and Monika Targan. A note on the oriented chromatic number of grids. *Inform. Process. Lett.*, 92(2):65–70, 2004.

[56] K. Thilagavathi and Vernold Vivin. J. Harmonious colouring of graphs. *Far East J. Math. Sci.*, 20(2):189–197, 2006.

[57] David R. Wood. Acyclic, star and oriented colourings of graph subdivisions. *Discrete Math. Theor. Comput. Sci.*, 7(1):37–50, 2005.

Departament de Matemàtica Aplicada II, Universitat Politècnica de Catalunya, Barcelona, Spain

E-mail address: david.wood@upc.edu