Transformation of Fractions into Simple Fractions in Divisive Meadows

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Abstract. Meadows are alternatives for fields with a purely equational axiomatization. At the basis of meadows lies the decision to make the multiplicative inverse operation total by imposing that the multiplicative inverse of zero is zero. Divisive meadows are meadows with the multiplicative inverse operation replaced by a division operation. Viewing a fraction as a term over the signature of divisive meadows that is of the form \( \frac{p}{q} \), we investigate which divisive meadows admit transformation of fractions into simple fractions, i.e. fractions without proper subterms that are fractions.

Keywords: divisive meadow, simple fraction, polynomial, rational number

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1 Introduction

To our knowledge, all existing definitions of a fraction are insufficiently precise to allow the validity of many non-trivial statements about fractions to be established. The work presented in this paper is concerned with the rigorous definition of a fraction and the validity of statements related to the question whether each fraction can be transformed into a simple fraction (colloquially described as a fraction where neither the numerator nor the denominator contains a fraction). This work is carried out in the setting of divisive meadows.

Because fields do not have a purely equational axiomatization, the axioms of a field cannot be used in applications of the theory of abstract data types to number systems based on rational, real or complex numbers. In [10], meadows are proposed as alternatives for fields with a purely equational axiomatization. At the basis of meadows lies the decision to make the multiplicative inverse operation total by imposing that the multiplicative inverse of zero is zero. A meadow is a commutative ring with a multiplicative identity element and a total multiplicative inverse operation satisfying the two equations \((x^{-1})^{-1} = x\) and \(x \cdot (x \cdot x^{-1}) = x\). It follows from the axioms of a meadow that the multiplicative inverse operation also satisfies the equation \(0^{-1} = 0\). All fields in which the multiplicative inverse of zero is zero, called zero-totalized fields, are meadows, but
not conversely. Because of their purely equational axiomatization, all meadows are total algebras and the class of all meadows is a variety.

In [6], divisive meadows are proposed. A divisive meadow is a commutative ring with a multiplicative identity element and a total division operation satisfying the three equations \(1 / (1 / x) = x\), \((x \cdot x) / x = x\), and \(x / y = x \cdot (1 / y)\). It follows from the axioms of a divisive meadow that the division operation also satisfies the equation \(x / 0 = 0\). We coined the alternative name inversive meadow for a meadow. The equational axiomatizations of inversive meadows and divisive meadows are essentially the same in the sense that they are definitionally equivalent.

We expect the zero-totalized multiplicative inverse and division operations of inversive and divisive meadows, which are conceptually and technically simpler than the conventional partial multiplicative inverse and division operations, to be useful in among other things mathematics education. We further believe that viewing fractions as terms over the signature of divisive meadows whose outermost operator is the division operator gives a rigorous definition of a fraction that can serve as a basis of a workable theory about fractions for teaching purposes at all levels of education (cf. [1]). Divisive meadows are more convenient than inversive meadows for the definition of fractions because, unlike the signature of inversive meadows, the signature of divisive meadows includes the division operator.

Viewing fractions as described above has two salient consequences: (i) fractions may contain variables and (ii) fractions may be interpreted in different divisive meadows. These consequences lead to the need to make a distinction between fractions and closed fractions and to consider properties of fractions relative to a particular divisive meadow. Viewing fractions as described above, many properties of fractions considered in the past turn out to be properties of closed fractions and/or properties of fractions relative to the divisive meadow of rational numbers.

For example, it is known from earlier work on meadows that closed fractions can be transformed into simple fractions, i.e. fractions without proper subterms that are fractions, if fractions are interpreted in the divisive meadow of rational numbers. Now the question arises whether the restriction to closed fractions can be dropped and whether this result goes through if fractions are interpreted in divisive meadows different from the divisive meadow of rational numbers.

In this paper, we investigate which divisive meadows admit transformation into simple fractions (for both the general case and the case of closed fractions). Some exemplary results are: (i) every model of the axioms of divisive meadow with a finite carrier admits transformation into simple fractions; (ii) every minimal model of the axioms of a divisive meadow with an infinite carrier does not admit transformation into simple fractions; (iii) the divisive meadow of rational numbers is the only minimal model of the axioms of a divisive meadow with an infinite carrier that admits transformation into simple fractions for closed fractions.
This paper is organized as follows. First, we give a survey of inversive meadows and divisive meadows which includes the signature and axioms for them, general results about them, and terminology used in the setting of meadows (Section 2). Next, we give the definitions concerning fractions and polynomials on which subsequent sections are based (Section 3) and establish some auxiliary results concerning divisive meadows which will be used in subsequent sections (Section 4). Then, we establish results about the transformation into simple fractions (Sections 5 and 6). Following this, we establish results that are related to the results in the preceding sections, but do not concern fractions (Section 7). Finally, we make some concluding remarks (Section 8).

We conclude this introduction with a corrective note on a remark made in [6] and later papers on meadows. Skew meadows, which differ from meadows only in that their multiplication is not required to be commutative, were already studied in [16,17], where they go by the name of desirable pseudo-fields. In 2009, we first read about desirable pseudo-fields in [19] and reported on it in [6]. However, we thought incorrectly at the time that desirable pseudo-fields were meadows. Hence, we accidentally mentioned that meadows were already introduced in [16].

2 Inversive Meadows and Divisive Meadows

In this section, we survey both inversive meadows and divisive meadows. Inversive meadows have been proposed as alternatives for fields with a purely equational axiomatization in [10]. Inversive meadows have been further investigated in e.g. [5,11,4,8,12] and applied in e.g. [9,2,7]. Divisive meadows, which are inversive meadows with the multiplicative inverse operation replaced by a division operation, have been proposed in [6]. In subsequent sections, only divisive meadows are needed. However, established results about inverse meadows will be used in proofs wherever it is justified by the definitional equivalence of the axiomatizations of inversive meadows and divisible meadows (see also the remarks following Theorem 4 below).

An inversive meadow is a commutative ring with a multiplicative identity element and a total multiplicative inverse operation satisfying two equations which imply that the multiplicative inverse of zero is zero. A divisive meadow is a commutative ring with a multiplicative identity element and a total division operation satisfying three equations which imply that division by zero always yields zero. Hence, the signature of both inversive and divisive meadows include the signature of a commutative ring with a multiplicative identity element.

The signature of commutative rings with a multiplicative identity element consists of the following constants and operators:

- the additive identity constant 0;
- the multiplicative identity constant 1;

1 An overview of all the work on meadows done to date and some open questions originating from that work can be found on [18].
Table 1. Axioms of a commutative ring with a multiplicative identity element

| Equation                                                                 | Revised Equation |
|-------------------------------------------------------------------------|------------------|
| \((x + y) + z = x + (y + z)\)                                           | \((x + y) \cdot z = x \cdot (y \cdot z)\) |
| \(x + y = y + x\)                                                      | \(x \cdot y = y \cdot x\) |
| \(x + 0 = x\)                                                          | \(x \cdot 1 = x\) |
| \(x + (-x) = 0\)                                                       | \(x \cdot (y + z) = x \cdot y + x \cdot z\) |

- the binary addition operator \(+\);
- the binary multiplication operator \(\cdot\);
- the unary additive inverse operator \(−\);

The signature of inversive meadows consists of the constants and operators from the signature of commutative rings with a multiplicative identity element and in addition:
- the unary multiplicative inverse operator \(\cdot\)^{-1}.

The signature of divisive meadows consists of the constants and operators from the signature of commutative rings with a multiplicative identity element and in addition:
- the binary division operator \(/\).

We write:

\[
\begin{align*}
\Sigma_{CR} & \text{ for } \{0, 1, +, ·, −\}, \\
\Sigma_{iMd} & \text{ for } \Sigma_{CR} \cup \{\cdot^{-1}\}, \\
\Sigma_{dMd} & \text{ for } \Sigma_{CR} \cup \{/\}.
\end{align*}
\]

We assume that there are infinitely many variables, including \(x\), \(y\) and \(z\). Terms are build as usual. We use infix notation for the binary operators, prefix notation for the unary operator \(−\), and postfix notation for the unary operator \(\cdot^{-1}\). We use the usual precedence convention to reduce the need for parentheses. Subtraction is introduced as an abbreviation as follows: \(p − q\) abbreviates \(p + (−q)\). We use the notation \(p^n\) for exponentiation with natural number exponents. For each term \(p\) over the signature of inversive meadows or the signature of divisive meadows and each natural number \(n\), the term \(p^n\) is defined by induction on \(n\) as follows: \(p^0 = 1\) and \(p^{n+1} = p^n \cdot p\). We use the notation \(\underline{n}\) for the numeral of \(n\). For each natural number \(n\), the term \(\underline{n}\) is defined by induction on \(n\) as follows: \(\underline{0} = 0\) and \(\underline{n + 1} = \underline{n} + 1\). For convenience, we extend the notation \(\underline{n}\) from natural numbers to integers by stipulating that \(−\underline{n} = \underline{-n}\).

The constants and operators from the signatures of inversive meadows and divisive meadows are adopted from rational arithmetic, which gives an appropriate intuition about these constants and operators.

A commutative ring with a multiplicative identity element is a total algebra over the signature \(\Sigma_{CR}\) that satisfies the equations given in Table 1. An inversive meadow is a total algebra over the signature \(\Sigma_{iMd}\) that satisfies the equations

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Table 2. Additional axioms for an inversive meadow

| Equation          |
|-------------------|
| $(x^{-1})^{-1} = x$ |
| $x \cdot (x \cdot x^{-1}) = x$ |

Table 3. Additional axioms for a divisive meadow

| Equation          |
|-------------------|
| $1 / (1 / x) = x$ |
| $(x \cdot x) / x = x$ |
| $x / y = x \cdot (1 / y)$ |

given in Tables 1 and 2. A *divisive meadow* is a total algebra over the signature $\Sigma^d_{Md}$ that satisfies the equations given in Tables 1 and 2.

We write:

- $E_{CR}$ for the set of all equations in Table 1,
- $E_{inv}$ for the set of all equations in Table 2,
- $E_{div}$ for the set of all equations in Table 3,
- $E^i_{Md}$ for $E_{CR} \cup E_{inv}$,
- $E^d_{Md}$ for $E_{CR} \cup E_{div}$.

Equations making the nature of the multiplicative inverse operation in inversive meadows more clear are derivable from the equations $E^i_{Md}$.

**Proposition 1.** The equations

$$0^{-1} = 0, \quad 1^{-1} = 1, \quad (-x)^{-1} = -(x^{-1}), \quad (x \cdot y)^{-1} = x^{-1} \cdot y^{-1}$$

are derivable from the equations $E^i_{Md}$.

**Proof.** Theorem 2.2 from [10] is concerned with the derivability of the first equation and Proposition 2.8 from [5] is concerned with the derivability of the last two equations. The derivability of the second equation is trivial.

The advantage of working with a total multiplicative inverse operation or a total division operation lies in the fact that conditions such as $x \neq 0$ in $x \neq 0 \Rightarrow x \cdot x^{-1} = 1$ are not needed to guarantee meaning.

An inversive or divisive meadow is *non-trivial* if it satisfies the *separation axiom*

$$0 \neq 1;$$

and it is an inversive or divisive *cancellation meadow* if it satisfies the *cancellation axiom*

$$x \neq 0 \land x \cdot y = x \cdot z \Rightarrow y = z.$$ 

In the case of an inversive meadow, the cancellation axiom is equivalent to the *general inverse law*

$$x \neq 0 \Rightarrow x \cdot x^{-1} = 1.$$
A totalized field is a total algebra over the signature $\Sigma_{Md}$ that satisfies the equations $E_{CR}$, the separation axiom, and the general inverse law. A zero-totalized field is a totalized field that satisfies in addition the equation $0^{-1} = 0$.

**Theorem 1.** The class of all non-trivial inversive cancellation meadows and the class of all zero-totalized fields are the same.

**Proof.** This is a corollary of Lemma 2.5 from [10]. □

Not all non-trivial inversive meadows are zero-totalized fields, e.g. the initial inversive meadow is not a zero-totalized field. Nevertheless, we have the following theorem.

**Theorem 2.** The equational theory of inversive meadows and the equational theory of zero-totalized fields are the same.

**Proof.** This is Theorem 3.10 from [5]. □

Theorem 2 can be read as follows: $E_{Md}^i$ is a finite basis for the equational theory of inversive cancellation meadows.

The inversive cancellation meadow that we are most interested in is $Q_0^i$, the zero-totalized field of rational numbers. $Q_0^i$ differs from the field of rational numbers only in that the multiplicative inverse of zero is zero.

**Theorem 3.** $Q_0^i$ is the initial algebra among the total algebras over the signature $\Sigma_{Md}^i$ that satisfy the equations

$$E_{Md}^i \cup \{ \frac{x}{y} = x \cdot y^{-1} \} \quad 2$$

or, equivalently, the equations

$$E_{Md}^i \cup \{ (1 + x^2 + y^2) \cdot (1 + x^2 + y^2)^{-1} = 1 \} .$$

**Proof.** As for the first set of equations, this is Theorem 3.1 in [10]. As for the second set of equations, this is Theorem 9 in [6]. □

The division operator can be explicitly defined in terms of the multiplicative inverse operator by the equation $x / y = x \cdot y^{-1}$ and the multiplicative inverse operator can be explicitly defined in terms of the division operator by the equation $x^{-1} = 1 / x$. In fact, $E_{Md}^i$ and $E_{Md}^d$ are essentially the same in the sense which is made precise in the following theorem.

**Theorem 4.** $E_{Md}^i$ is definitionally equivalent to $E_{Md}^d$ 3 i.e.

$$E_{Md}^i \cup \{ x / y = x \cdot y^{-1} \} \vdash E_{Md}^d \cup \{ x^{-1} = 1 / x \}$$

and

$$E_{Md}^d \cup \{ x^{-1} = 1 / x \} \vdash E_{Md}^i \cup \{ x / y = x \cdot y^{-1} \} .$$

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2 We write $\mathbb{N}^+$ for the set \{ $n \in \mathbb{N} \mid n \neq 0$ \} of positive natural numbers.

3 The notion of definitional equivalence originates from [13], where it was introduced, in the setting of first-order theories, under the name of synonymy. In [21], the notion of definitional equivalence was introduced in the setting of equational theories under the ambiguous name of equivalence. An abridged version of [21] appears in [14].
Proof. Because \( E_{\text{inv}} \) and \( E_{\text{div}} \) have \( E_{\text{CR}} \) in common, in one direction, we only have to prove the derivability of \( E_{\text{div}} \cup \{ x^{-1} = 1 / x \} \) and, in the other direction, we only have to prove the derivability of \( E_{\text{inv}} \cup \{ x / y = x \cdot y^{-1} \} \). The derivability of all equations involved is trivial. \( \square \)

By the definitional equivalence of \( E_{\text{Md}}^1 \) and \( E_{\text{Md}}^3 \), we have:

(a) there exist a mapping \( \epsilon \) from the set of all equations over \( \Sigma_{\text{Md}}^1 \) to the set of all equations over \( \Sigma_{\text{Md}}^3 \) and a mapping \( \epsilon' \) from the set of all equations over \( \Sigma_{\text{Md}}^1 \) to the set of all equations over \( \Sigma_{\text{Md}}^3 \) such that, for each equation \( \phi \) over \( \Sigma_{\text{Md}}^1 \) and each equation \( \phi' \) over \( \Sigma_{\text{Md}}^3 \),

\[
E_{\text{Md}}^1 \vdash \epsilon(\phi) \quad \text{if} \quad E_{\text{Md}}^3 \vdash \phi,
\]

\[
E_{\text{Md}}^3 \vdash \epsilon'(\phi) \quad \text{if} \quad E_{\text{Md}}^1 \vdash \phi'.
\]

(b) there exist a mapping \( \alpha \) from the class of all divisive meadows to the class of all inversive meadows that maps each divisive meadow \( \mathcal{M}' \) to the restriction to \( \Sigma_{\text{Md}}^1 \) of the unique expansion of \( \mathcal{M}' \) for which \( E_{\text{Md}}^3 \cup \{ x^{-1} = 1 / x \} \) holds and a mapping \( \alpha' \) from the class of all inversive meadows to the class of all divisive meadows that maps each inversive meadow \( \mathcal{M} \) to the restriction to \( \Sigma_{\text{Md}}^1 \) of the unique expansion of \( \mathcal{M} \) for which \( E_{\text{Md}}^1 \cup \{ x / y = x \cdot y^{-1} \} \) holds such that \( \alpha \circ \alpha' \) and \( \alpha' \circ \alpha \) are identity mappings.

Let \( \epsilon \) and \( \alpha \) be as under (a) and (b) above. Then it follows that, for each equation \( \phi \) over \( \Sigma_{\text{Md}}^1 \) and divisive meadow \( \mathcal{M} \):

- \( E_{\text{Md}}^1 \vdash \phi \) iff \( E_{\text{Md}}^1 \vdash \epsilon(\phi) \);
- \( \mathcal{M} \models \phi \) iff \( \alpha(\mathcal{M}) \models \epsilon(\phi) \).

From many results about inversive meadows (including the ones presented above), counterparts about divisive meadows follow immediately using these consequences of the definitional equivalence of \( E_{\text{Md}}^1 \) and \( E_{\text{Md}}^3 \). This is the main reason why the survey given in this section is not restricted to divisive meadows.

Henceforth, “meadow” without “inversive” or “divisive” as qualifier stands for “inversive or divisive meadow”.

A meadow is finite if its carrier is finite and it is infinite if its carrier is infinite. A meadow \( \mathcal{M} \) is a (proper) submeadow of a meadow \( \mathcal{M}' \) if \( \mathcal{M} \) is a (proper) subalgebra of \( \mathcal{M}' \). A meadow is minimal if it does not have a proper submeadow. The characteristic of a meadow is the smallest \( k \in \mathbb{N}^+ \) for which it satisfies \( k = 0 \). A meadow is said to have characteristic 0 if there does not exist an \( k \in \mathbb{N}^+ \) for which it satisfies \( k = 0 \). A \( k \in \mathbb{N}^+ \) is called square-free if it is the product of distinct prime numbers.

If a meadow has characteristic 0, then it also satisfies, for all \( k, k' \in \mathbb{N}, k \neq k' \) if \( k \neq k' \). Hence, a meadow has characteristic 0 only if its minimal submeadow is infinite.

The infinite divisive meadow that we are most interested in is \( \mathbb{Q}^d_0 \), the zero-totalized field of rational numbers with the multiplicative inverse operation replaced by a division operation. It follows immediately from Theorems 3 and 4.
that $Q^d_0$ is the initial algebra among the total algebras over the signature $\Sigma^d_{Md}$ that satisfy the equations $E^d_{Md} \cup \{n \div n = 1 \mid n \in \mathbb{N}^+\}$.

The finite divisive meadows that we are most interested in are, for each square-free $k \in \mathbb{N}^+$, $M^d_k$, the initial algebra among the total algebras over the signature $\Sigma^d_{Md}$ that satisfy the equations $E^d_{Md} \cup \{k = 0\}$. It follows immediately from Lemma 4.6 in [5] and Theorem 4 in this paper that each minimal divisive meadow of characteristic $k$ is isomorphic to $M^d_k$. It follows immediately from Theorem 4.4 and Lemma 4.6 in [5], Corollary 3.10 in [12], and Theorem 4 in this paper that, if $k$ is a prime number, $M^d_k$ is the zero-totalized prime field of characteristic $k$ with the multiplicative inverse operation replaced by a division operation.

Several results established in Sections 4–6 are exclusively concerned with the divisive meadow $Q^d_0$. The divisive meadows $M^d_k$ play a role in various proofs given in Sections 5–7.

3 Definitions Concerning Fractions and Polynomials

In this section, we give the definitions concerning fractions on which the subsequent sections are based. Because polynomials play a role in those sections as well, we also give several definitions concerning polynomials in the setting of divisive meadows.

Henceforth, we will use the following convenient notational convention. If we introduce a term $t$ as $t(x_1, \ldots, x_n)$, where $x_1, \ldots, x_n$ are distinct variables, this indicates that all variables that have occurrences in $t$ are among $x_1, \ldots, x_n$. In the same context, $t(t_1, \ldots, t_n)$ is the term obtained by simultaneously replacing in $t$ all occurrences of $x_1$ by $t_1$ and \ldots and all occurrences of $x_n$ by $t_n$.

Fractions are viewed as terms over the signature of divisive meadows that are of a particular form. This means that fractions may contain variables and may be interpreted in different divisive meadows. Thus, the view of fractions as terms leads to the need to make a distinction between fractions and closed fractions and to consider properties of fractions relative to a particular divisive meadow. In this light, the definitions given below speak for themselves.

The following four definitions concern fractions by themselves:

- a fraction is a term over the signature $\Sigma^d_{Md}$ whose outermost operator is the operator $/$;
- a simple fraction is a fraction of which no proper subterm is a fraction;
- a closed fraction is a fraction that is a closed term;
- a simple closed fraction is a simple fraction that is a closed term.

Let $E \supseteq E^d_{Md}$ be a set of equations over the signature $\Sigma^d_{Md}$, and let $M$ be a model of $E$. The following four definitions concern the transformation of terms into simple fractions:

- $M$ admits transformation into simple fractions if, for each term $p$ over the signature $\Sigma^d_{Md}$, there exists a simple fraction $q$ such that $M \models p = q$.
- \( \mathcal{M} \) admits transformation into simple fractions for closed terms if, for each closed term \( p \) over the signature \( \Sigma_{\text{Md}}^d \), there exists a simple closed fraction \( q \) such that \( \mathcal{M} \models p = q \);
- \( E \) admits transformation into simple fractions if, for each term \( p \) over the signature \( \Sigma_{\text{Md}}^d \), there exists a simple fraction \( q \) such that \( E \vdash p = q \);
- \( E \) admits transformation into simple fractions for closed terms if, for each closed term \( p \) over the signature \( \Sigma_{\text{Md}}^d \), there exists a simple closed fraction \( q \) such that \( E \vdash p = q \).

We believe that a fraction corresponds to what is usually meant by an algebraic fraction. We are not sure to what extent a closed fraction corresponds to what is usually meant by a numerical fraction because definitions of the latter are mostly so imprecise that form and meaning seem to be mixed up.

Notice that, for each term \( p \) over the signature \( \Sigma_{\text{Md}}^d \), \( E_{\text{Md}}^d \vdash p = p / 1 \). This means that each term \( p \) over the signature \( \Sigma_{\text{Md}}^d \) that is not a fraction can be turned into a fraction in a trivial way.

In subsequent sections, we will phrase some results about the transformation of terms into simple fractions in which we refer in one way or another to polynomials. Therefore, we also give several definitions concerning polynomials in the setting of divisive meadows. A polynomial as defined below corresponds to what is usually meant by a (univariate) polynomial.

Let \( x, y \) be variables, and let \( \mathcal{M} \) be a divisive meadow. The following definitions concern polynomials:

- a polynomial in the variable \( x \) is a term over the signature \( \Sigma_{\text{Md}}^d \) in which the operator / does not occur and variables other than \( x \) do not occur;
- a root of a polynomial \( f(x) \) over \( \mathcal{M} \) is an element \( v \) of the carrier of \( \mathcal{M} \) such that \( \mathcal{M} \) satisfies the equation \( f(x) = 0 \) if the value assigned to \( x \) is \( v \);
- the polynomial function induced by a polynomial \( f(x) \) over \( \mathcal{M} \) is the unary function \( F \) on the carrier of \( \mathcal{M} \) such that, for each element \( v \) of the carrier of \( \mathcal{M} \), \( F(v) \) is the interpretation of \( f(x) \) in \( \mathcal{M} \) if the value assigned to \( x \) is \( v \);
- a polynomial \( f(x) \) is \( \mathcal{M} \)-equivalent to a polynomial \( g(y) \), written \( f(x) \equiv_{\mathcal{M}} g(y) \), if \( \mathcal{M} \models f(x) = g(x) \);
- a polynomial is in canonical form if it is of the form \( a_n \cdot x^n + \ldots + a_1 \cdot x + a_0 \), where \( a_i \in \{ m \mid m \in \mathbb{Z} \} \) for each \( i \in \{1, \ldots, n\} \).

The closed terms \( a_i \in \{ m \mid m \in \mathbb{Z} \} \) occurring in a polynomial in canonical form are called coefficients.

It is a generally known fact that, for each polynomial \( f(x) \), there exists a polynomial \( g(x) \) in canonical form such that \( E_{\text{CR}} \vdash f(x) = g(x) \). This fact can be straightforwardly proved by induction on the structure of a polynomial. The following definitions concerning polynomials are based on this fact:

- a polynomial \( f(x) \) is non-trivial over \( \mathcal{M} \) if there exist a polynomial \( g(x) \) in canonical form, an \( a \in \{ m \mid m \in \mathbb{Z} \} \), and an \( i \in \mathbb{N} \) such that \( E_{\text{CR}} \vdash f(x) = g(x), a \cdot x^i \) is a summand of \( g(x) \), and \( \mathcal{M} \not\models a = 0 \);
- a polynomial \( f(x) \) is constant over \( \mathcal{M} \) if \( f(x) \) is non-trivial over \( \mathcal{M} \) and there exists an \( a \in \{ m \mid m \in \mathbb{Z} \} \) such that \( \mathcal{M} \models f(x) = a \);
the degree of a polynomial $f(x)$ over $\mathcal{M}$ is:

- if $f(x)$ is non-trivial and not constant over $\mathcal{M}$, then the largest $i \in \mathbb{N}^+$ for which there exists a polynomial $g(x)$ in canonical form and an $a \in \{m \mid m \in \mathbb{Z}\}$ such that $E_{\text{CR}} \vdash f(x) = g(x)$, $a \cdot x^i$ is a summand of $g(x)$, and $\mathcal{M} \not| a = 0$;
- if $f(x)$ is non-trivial and constant over $\mathcal{M}$, then 0;
- if $f(x)$ is not non-trivial over $\mathcal{M}$, then undefined.

Notice that a simple fraction is a fraction of which the two outermost proper subterms are polynomials.

Henceforth, we will often leave out the qualifier “over $\mathcal{M}$” used in the definienda above in contexts where only one divisive meadow is under discussion.

## 4 Auxiliary Results Concerning Divisive Meadows

In this section, we establish some results concerning divisive meadows that will be used in Sections 5 and 6 to establish results about the transformation of terms into simple fractions in divisive meadows.

The following proposition, which concerns equations making the nature of the division operation in divisive meadows more clear, is useful in many proofs.

**Proposition 2.** The equations

\[
\frac{1}{0} = 0, \quad \frac{1}{1} = 1, \quad \frac{1}{(-x)} = -(1/x), \quad 1/(x \cdot y) = (1/x) \cdot (1/y), \\
(x/y) \cdot (z/w) = (x \cdot z)/(y \cdot w), \quad (x/y)/(z/w) = (x \cdot w)/(y \cdot z)
\]

are derivable from the equations $E_{\text{Md}}^d$.

**Proof.** The derivability of the first four equations follow immediately from Proposition 1 and the definitional equivalence of $E_{\text{Md}}^i$ and $E_{\text{Md}}^d$. The last two equations are easily derivable using the fourth equation. $\square$

The following proposition is also useful in several proofs.

**Proposition 3.** For all $n, m \in \mathbb{N}$, the equations $n + m = m + n$ and $n \cdot m = m \cdot n$ are derivable from the equations $E_{\text{Md}}^d$.

**Proof.** This is proved like Lemma 1 in [6]. $\square$

Closed terms over the signature of divisive meadows can be reduced to a basic term. The set $\mathcal{B}$ of basic terms over $\Sigma_{\text{Md}}^d$ is inductively defined by the following rules:

- $0 \in \mathcal{B}$;
- if $n, m \in \mathbb{N}^+$, then $n/m \in \mathcal{B}$;
- if $n, m \in \mathbb{N}^+$, then $-(n/m) \in \mathcal{B}$;
- if $p, q \in \mathcal{B}$, then $p + q \in \mathcal{B}$. 

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Theorem 5. For all closed terms $p$ over $\Sigma^{\text{Md}}_d$, there exists a $q \in B$ such that $E^{\text{Md}}_d \vdash p = q$.

Proof. The proof is straightforward by induction on the structure of $p$. If $p$ is of the form $0$, $1$, $p' + q'$, or $p' \cdot q'$, then it is trivial to show that there exists a $q \in B$ such that $E^{\text{Md}}_d \vdash p = q$. If $p$ is of the form $-p'$, $p' \cdot q'$ or $p' / q'$, then it follows immediately from the induction hypothesis and the following claims:

- for all $p' \in B$, there exists a $p'' \in B$ such that $E^{\text{Md}}_d \vdash -p' = p''$;
- for all $p', q' \in B$, there exists a $p'' \in B$ such that $E^{\text{Md}}_d \vdash p' \cdot q' = p''$;
- for all $p', q' \in B$, there exists a $p'' \in B$ such that $E^{\text{Md}}_d \vdash p' / q' = p''$.

These claims are easily proved by induction on the structure of $p'$, using Proposition 3. □

It is well known that closed terms over $\Sigma^{\text{Md}}_d$ in which the operator $/$ does not occur can be reduced to a closed term of the form $0$, $n$ or $-n$.

Proposition 4. For all closed terms $p$ over $\Sigma^{\text{CR}}_d$, $E^{\text{Md}}_d \vdash p = 0$ or there exists an $n \in \mathbb{N}^+$ such that $E^{\text{Md}}_d \vdash p = n$ or $E^{\text{Md}}_d \vdash p = -n$.

Proof. The proof is along the lines of the proof of Theorem 5 but simpler. □

Each infinite minimal divisive meadow has $\mathbb{Q}_0^d$ as a homomorphic image.  

Theorem 6. $\mathbb{Q}_0^d$ is a homomorphic image of each infinite minimal divisive meadow.

Proof. In order to prove this theorem by contradiction, assume that $\mathcal{M}$ is an infinite minimal divisive meadow that does not have $\mathbb{Q}_0^d$ as a homomorphic image. Then there exists a closed equation $p = q$ such that $\mathcal{M} \models p = q$ and $\mathbb{Q}_0^d \not\models p = q$. Because $E^{\text{Md}}_d \cup \{p = q\} \vdash p - q = 0$, there also exists a closed equation $p' = 0$ such that $\mathcal{M} \models p' = 0$ and $\mathbb{Q}_0^d \not\models p' = 0$. Because of Theorem 5, we may assume that the closed term $p'$ in such an equation is a basic term.

Let $p \in B$ be such that $\mathcal{M} \models p = 0$ and $\mathbb{Q}_0^d \not\models p = 0$. Then there exists an $n \in \mathbb{N}^+$ such that $E^{\text{Md}}_d \cup \{p = 0\} \vdash n = 0$. We prove this by induction on the number of different subterms of $p$ of the form $p' + q'$ in which $p'$ or $q'$ has a subterm of the form $k / l$ with $k, l \in \mathbb{N}^+$ and $l \neq 1$. The basis step is easily proved: $E^{\text{Md}}_d \cup \{p = 0\} \vdash n = 0$ follows immediately from the fact that $p$ is of the form $k / l$ with $k \in \mathbb{N}^+$. The inductive step is proved in the following way. Necessarily, $p$ is of the form $C[q + q']$ where $q = k / l$ or $q = -(k / l)$ for some $k, l \in \mathbb{N}^+$, and $q' = k'/ l'$ or $q' = -(k'/ l')$ for some $k', l' \in \mathbb{N}^+$. We only consider the case that $q = k / l$ and $q' = k'/ l'$. The other cases are proved analogously. We have $E^{\text{Md}}_d \cup \{p = 0\} \vdash (k / l) \cdot (l' / l') \cdot C[k / l + k' / l'] = 0$. It is easily proved that $E^{\text{Md}}_d \vdash (r / r) \cdot C[s] = (r / r) \cdot C[(r / r) \cdot s]$ for all terms $r$ and $s$ and contexts $C[\cdot]$ over $\Sigma^{\text{Md}}_d$ (cf. the proof of Corollary 3.1 in [4]). Hence, $E^{\text{Md}}_d \cup \{p = 0\} \vdash (k / l) \cdot (l' / l') \cdot C[(l' / l') \cdot (k / l) + (k / l) \cdot (k' / l')] = 0$ and $\mathbb{Q}_0^d$ is the divisive meadow of rational numbers introduced at the end of Section 2.
so $E_{\Sigma_{\text{Md}}}^d \cup \{ p = 0 \} \vdash (\ell / \ell) \cdot (\ell' / \ell') \cdot C((k \cdot \ell + k' \cdot \ell) / (\ell \cdot \ell')) = 0$. From this, it follows that $\mathcal{M} \models (\ell / \ell) \cdot (\ell / \ell') \cdot C((k \cdot \ell + k' \cdot \ell) / (\ell \cdot \ell')) = 0$, whereas $Q_0^d \not\models (\ell / \ell) \cdot (\ell / \ell') \cdot C((k \cdot \ell + k' \cdot \ell) / (\ell \cdot \ell')) = 0$ since $Q_0^d \not\models p = 0$, $Q_0^d \models \ell / \ell = 1$, and $Q_0^d \models \ell' / \ell' = 1$. In the term $(\ell / \ell) \cdot (\ell / \ell') \cdot C((k \cdot \ell + k' \cdot \ell) / (\ell \cdot \ell'))$, the number used for the induction is one less than in $p$. This means that we can apply the induction hypothesis and from that it follows that there exists an $n \in \mathbb{N}^+$ such that $E_{\Sigma_{\text{Md}}}^d \cup \{ \ell / \ell \} \vdash \frac{p}{n} = 0$.

Let $n \in \mathbb{N}^+$ be such that $E_{\Sigma_{\text{Md}}}^d \cup \{ p = 0 \} \vdash \frac{n}{n} = 0$. Then it follows immediately that $E_{\Sigma_{\text{Md}}}^d \cup \{ p = 0 \} \vdash \frac{m}{m} = 0$, where $m$ is the product of the prime factors of $n$ (each with multiplicity 1). Hence, each model of $E_{\Sigma_{\text{Md}}}^d \cup \{ p = 0 \}$ has a non-zero square-free characteristic. Now it is an immediate corollary of Theorem 4.4 and Lemma 4.6 in [5] that each minimal inversive meadow of non-zero square-free characteristic is finite. This finiteness result carries over to minimal divisive meadows of non-zero square-free characteristic by Theorem [4]. Consequently, each minimal model of $E_{\Sigma_{\text{Md}}}^d \cup \{ p = 0 \}$ is finite. This contradicts the assumed infinity of $\mathcal{M}$. Hence, there does not exist an infinite minimal divisive meadow that does not have $Q_0^d$ as a homomorphic image. 

\section{Transformation of Fractions, the General Case}

In this section, we establish results about the transformation into simple fractions for terms over the signature of divisive meadows. The results concerned are not restricted to closed terms. In Section 5, we will establish results that are restricted to closed terms.

The first result concerns finite divisive meadow.

\begin{theorem}
Every finite divisive meadow admits transformation into simple fractions.
\end{theorem}

\begin{proof}
Let $\mathcal{M}$ be a finite divisive meadow. Then there exist $n, m \in \mathbb{N}$ such that $\mathcal{M} \models x^n = x^m$. Then $E_{\Sigma_{\text{Md}}}^d \cup \{ x^n = x^m \} \vdash \frac{1}{x} = x^{2(n-m)-1}$. We prove this using that, for each $k \in \mathbb{N}$, $E_{\Sigma_{\text{Md}}}^d \vdash \frac{x}{x} = x^{k+1}/x^k$ and $E_{\Sigma_{\text{Md}}}^d \vdash \frac{1}{x} = x^k/x^{k+1}$ as follows: $1/x = x^n/x^{n+1} = x^{2n-m}/x^{m+1} = x^{2(n-m)-1}$.

From $\mathcal{M} \models x^n = x^m$ and $E_{\Sigma_{\text{Md}}}^d \cup \{ x^n = x^m \} \vdash \frac{1}{x} = x^{2(n-m)-1}$, it follows that $\mathcal{M} \models \frac{1}{x} = x^{2(n-m)-1}$. Consequently, $\mathcal{M}$ admits that each term $p$ over the signature $\Sigma_{\text{Md}}^d$ is transformed into a simple fraction.
\end{proof}

The next two results tell us that some but not all infinite divisive meadows admit transformation into simple fractions.

\begin{theorem}
There exists an infinite divisive meadow that admits transformation into simple fractions.
\end{theorem}

\begin{proof}
Let $c_i$ be a constant for each $i \in \mathbb{N}$, let $\Sigma = \Sigma_{\text{Md}}^d \cup \{ c_i \mid i \in \mathbb{N} \}$, let $E = E_{\Sigma_{\text{Md}}}^d \cup \{ 0, x^2 = x \}$, and let $\mathcal{I}$ be the initial algebra among the total algebras over the signature $\Sigma$ that satisfy the equations $E$. Then $\mathcal{I} \models \frac{1}{x} = x$.

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We prove this using that $E_{Md} \vdash x = x^2 / x$ and $E_{Md} \vdash 1 / x = x / x^2$ as follows:

$$\frac{1}{x} = \frac{x}{x^2} = \frac{x}{x} = x$$

Because $I \models 1 / x = x$, $I$ admits that each term $p$ over the signature $\Sigma$ is transformed into a simple fraction. Moreover, $I$ is an infinite divisive meadow.

Take two arbitrary constants $c_i$ and $c_j$ with $i \neq j$. We find a total algebra over the signature $\Sigma$ that satisfies the equations $E$ but not the equation $c_i = c_j$ by taking the interpretation of $0$ as the interpretation of $c_i$, taking the interpretation of $1$ as the interpretation of $c_j$, and taking the interpretation of either $0$ or $1$ as the interpretation of the other constants. If follows from the existence of such an algebra that $E \not\vdash c_i = c_j$. Consequently, $I \not\models c_i = c_j$. Hence, $I$ is an infinite divisive meadow.

**Theorem 9.** $Q_0^d$ does not admit transformation into simple fractions.

**Proof.** In order to prove this theorem by contradiction, assume that $Q_0^d$ admits transformation into simple fractions. Then there exist polynomials $f(x)$ and $g(x)$ such that $Q_0^d \models 1 + 1 / x = f(x) / g(x)$.

Let $f(x)$ and $g(x)$ be polynomials such that $Q_0^d \models 1 + 1 / x = f(x) / g(x)$. Then $Q_0^d \models g(0) \neq 0$ because $Q_0^d \models 1 = f(0) / g(0)$. From this, it follows that $f(x) / g(x)$, interpreted as a real function, is continuous on a closed interval $[0, \epsilon]$ for $\epsilon > 0$ so small that no root of $g(x)$ is in this interval. Now let $a$ be the maximal absolute value of $f(x)$ on the interval $[0, \epsilon]$ and let $b$ be the minimal absolute value of $g(x)$ on this interval. Then $f(x) / g(x) \leq a / b$ on the interval $[0, \epsilon]$.

Clearly, there exists a positive rational number $q$ in the interval $(0, \epsilon)$ so small that $1 + 1 / q > a / b$. Let $q$ be a positive rational number $q$ such that $1 + 1 / q > a / b$. Then also $f(q) / g(q) > a / b$ because of the assumption that $Q_0^d \models 1 + 1 / x = f(x) / g(x)$. This contradicts the fact that $f(x) / g(x) \leq a / b$ on the interval $(0, \epsilon)$. Hence, $Q_0^d$ does not admit transformation into simple fractions. 

Theorems 6 and 9 give rise to several corollaries.

**Corollary 1.** A divisive meadow of characteristic 0 does not admit transformation into simple fractions.

**Proof.** This follows immediately from Theorems 6 and 9, and the fact that a meadow has characteristic 0 only if its minimal submeadow is infinite (see the remark immediately after the definition of characteristic in Section 2).

**Corollary 2.** A minimal divisive meadow admits transformation into simple fractions if and only if it is finite.

**Proof.** This follows immediately from Theorems 6, 7, and 9.

The next result is one in which admitting transformation into simple fractions is related to the existence of a polynomial with a particular property.
Theorem 10. A divisive meadow \( \mathcal{M} \) admits transformation into simple fractions only if there exists a non-trivial polynomial \( f(x) \) such that each element of the carrier of \( \mathcal{M} \) is a root of \( f(x) \).

Proof. Let \( \mathcal{M} \) be a divisive meadow that admits transformation into simple fractions. Then there exist polynomials \( f(x) \) and \( g(x) \) such that \( \mathcal{M} \models 1 + 1/x = f(x)/g(x) \).

Let \( f(x) \) and \( g(x) \) be polynomials such that \( \mathcal{M} \models 1 + 1/x = f(x)/g(x) \). Substitution of \( x \) by 0 yields \( \mathcal{M} \models f(0)/g(0) \), so \( \mathcal{M} \not\models g(0) = 0 \). Multiplication by \( x^2 \cdot g(x)^2 \) yields \( \mathcal{M} \models x^2 \cdot g(x)^2 + (1/x) \cdot x^2 \cdot g(x)^2 = (f(x)/g(x)) \cdot x^2 \cdot g(x)^2 \), and then applying the axiom \( (x \cdot x)/x = x \) twice yields \( \mathcal{M} \models x^2 \cdot g(x)^2 + x \cdot g(x)^2 = f(x) \cdot x^2 \cdot g(x) \). Hence, \( \mathcal{M} \models x^2 \cdot g(x)^2 + x \cdot g(x)^2 - f(x) \cdot x^2 \cdot g(x) = 0 \). This means that each element of the carrier of \( \mathcal{M} \) is the root of the polynomial \( x^2 \cdot g(x)^2 + x \cdot g(x)^2 - f(x) \cdot x^2 \cdot g(x) \). It remains to be proved that this polynomial is non-trivial.

Considering that \( E_{CR} \vdash x^2 \cdot g(x)^2 + x \cdot g(x)^2 - f(x) \cdot x^2 \cdot g(x) = g(x)^2 \cdot x + (g(x)^2 - f(x) \cdot g(x)) \cdot x^2 \), it must be the case that a canonical form of \( x^2 \cdot g(x)^2 + x \cdot g(x)^2 - f(x) \cdot x^2 \cdot g(x) \) has the constant term 0 and the linear term \( g(0)^2 \cdot x \) as summands. From \( \mathcal{M} \not\models g(0) = 0 \), it follows that \( \mathcal{M} \not\models g(0)^2 = 0 \). Hence, \( x^2 \cdot g(x)^2 + x \cdot g(x)^2 - f(x) \cdot x^2 \cdot g(x) \) is non-trivial.

Theorem 10 gives rise to several corollaries.

Corollary 3. A divisive meadow whose carrier contains an element that is not the root of a non-trivial polynomial does not admit transformation into simple fractions.

Proof. This follows immediately from Theorem 10.

Corollary 4. A divisive meadow admits transformation into simple fractions only if there exists an \( n \in \mathbb{N} \) such that each element of its carrier is the root of a non-trivial polynomial of degree \( n \) or less.

Proof. This follows immediately from Theorem 10.

The next result gives us sufficient and necessary conditions of admitting transformation into simple fractions for divisive meadows of prime characteristic.

Theorem 11. Let \( k \in \mathbb{N} \) be a prime number, and let \( \mathcal{M} \) be a divisive meadow of characteristic \( k \). Then the following are equivalent:

(i) \( \mathcal{M} \) admits transformation into simple fractions;
(ii) there exists an \( n \in \mathbb{N} \) such that each element of the carrier of \( \mathcal{M} \) is the root of a non-trivial polynomial of degree \( n \) or less;
(iii) there exists a non-trivial polynomial \( f(x) \) such that each element of the carrier of \( \mathcal{M} \) is a root of \( f(x) \);
(iv) there exists an \( n \in \mathbb{N}^+ \) such that \( \mathcal{M} \models 1/x = x^n \).

Proof. Assume (i). Then (ii) follows immediately from Corollary 4.

Assume (ii). Let \( n \in \mathbb{N} \) be such that each element of the carrier of \( \mathcal{M} \) is the root of a non-trivial polynomial of degree \( n \) or less. We know from Theorem 4.4
and Lemma 4.6 in \[5\] and Theorem 4 in this paper that a minimal divisive meadow of prime characteristic is a finite divisive meadow. Because \(M\) is of prime characteristic and the interpretation of coefficients in \(M\) is the same as the interpretation of coefficients in the minimal divisive submeadow of \(M\), the set of non-trivial polynomials of degree \(n\) or less modulo \(M\)-equivalence is finite. Let \(f(x)\) be the product of the elements of a transversal for this set. Then \(f(x)\) is a non-trivial polynomial and each element of the carrier of \(M\) is a root of \(f(x)\).

Assume (iii). Let \(f(x)\) be a non-trivial polynomial such that each element of the carrier of \(M\) is a root of \(f(x)\). Then \(M \models f(x) = 0\). Assume that 
\[
 f(x) = a_n \cdot x^n + \ldots + a_1 \cdot x + a_0 \quad \text{with} \quad M \models a_n \neq 0.
\]
We know from Theorem 4.4 and Lemma 4.6 in \[5\], Corollary 3.10 in \[12\], and Theorem 4 in this paper that a minimal divisive meadow of prime characteristic is a finite divisive cancellation meadow. Because \(M\) is of prime characteristic, the interpretation of coefficients in \(M\) is the same as the interpretation of coefficients in the minimal divisive submeadow of \(M\), and \(M \models a_n \neq 0\), we have \(M \models a_n / a_n = 1\). Dividing both sides of the equation \(f(x) = 0\) by \(a_n\) yields \(M \models x^n + \ldots + (a_1 / a_n) \cdot x + a_0 / a_n = 0\). From this it follows by induction on \(i\) that, for all \(i \geq n\), there exists a polynomial \(g(x)\) of degree less than \(n\) such that \(M \models x^i = g(x)\). From this and the fact that the set of polynomials of degree \(n - 1\) or less modulo \(M\)-equivalence is finite (as explained above), the polynomials from the sequence \(x, x^2, x^3, \ldots\) cannot be all different modulo \(M\)-equivalence. In other words, there exist \(l, m \in \mathbb{N}^+\) with \(l > m\) such that \(M \models x^l = x^m\). Let \(l, m \in \mathbb{N}^+\) be such that \(l > m\) and \(M \models x^l = x^m\). Then we can prove like in the proof of Theorem 7 that \(M \models 1 / x = x^{2(l-m)-1}\).

Assume (iv). Let \(n \in \mathbb{N}^+\) be such that \(M \models 1 / x = x^n\). Then it follows immediately that, for each term \(p\) over \(\Sigma^d_{Md}\), there exists a term \(q\) over \(\Sigma^d_{Md}\) of which no subterm is a fraction such that \(M \models p = q\). Hence, \(M\) admits transformation into simple fractions.

The next theorem tells us that admitting transformation into simple fractions is a property of divisive meadows that cannot be expressed as a first-order theory.

**Theorem 12.** Admitting transformation into simple fractions is not an elementary property of divisive meadows.

*Proof.* In order to prove this theorem by contradiction, assume that there exists a first-order theory \(T\) over \(\Sigma^d_{Md}\) such that the models of \(E^d_{Md} \cup T\) are precisely the divisive meadows with the mentioned property.

Let \(T\) be a first-order theory over \(\Sigma^d_{Md}\) such that the models of \(E^d_{Md} \cup T\) are precisely the divisive meadows with the mentioned property. Because of Theorem 4 for each \(k' \in \mathbb{N}\), there exists a square-free \(k'' \in \mathbb{N}\) with \(k'' > k'\) such that \(Md^d_{k''} \models E^d_{Md} \cup T \cup \{k \neq 0 \mid k \in \mathbb{N}^+, k \leq k'\}\)\footnote{\(Md^d_k\), for square-free \(k \in \mathbb{N}^+\), was introduced at the end of Section 2}. From this, it follows that, for each \(k' \in \mathbb{N}\), there exists a divisive meadow \(M\) such that \(M \models E^d_{Md} \cup T \cup \{k \neq 0 \mid k \in \mathbb{N}^+, k \leq k'\}\). Hence, by the compactness of first-order logic, there
exists a divisive meadow $\mathcal{M}'$ such that $\mathcal{M}' \models E^d_{\text{Md}} \cup T \cup \{ k \neq 0 \mid k \in \mathbb{N}^+\}$.

In other words, there exists a divisive meadow $\mathcal{M}$ of characteristic 0 such that $\mathcal{M} \models E^d_{\text{Md}} \cup T$. This contradicts Corollary 1. Hence, there does not exist a first-order theory $T$ over $\Sigma^d_{\text{Md}}$ such that the models of $E^d_{\text{Md}} \cup T$ are precisely the divisive meadows with the mentioned property. \qed

The next theorem tells us that each divisive meadow admits transformation into sums of simple fractions.

**Theorem 13.** For each term $p$ over the signature $\Sigma^d_{\text{Md}}$, there exists a finite number of simple fractions, say $q_1, \ldots, q_n$, such that $E^d_{\text{Md}} \models p = q_1 + \ldots + q_n$.

**Proof.** By Theorem 2.1 in [4], for each term $p$ over the signature $\Sigma^d_{\text{Md}}$, there exists a term $q$ in standard meadow form such that $E^d_{\text{Md}} \models p = q$. By the distributivity of multiplication over addition, each term in standard meadow form is derivably equal to a sum of terms in standard meadow form of level 0; and terms in standard meadow form of level 0 are of the form $p' \cdot q'^{-1}$ with $p'$ and $q'$ polynomials. By Theorem 1, this proves the current theorem. \qed

A question arising from Theorem 13 is whether there exists a natural number $k$ such that each term over the signature $\Sigma^d_{\text{Md}}$ is derivably equal to a sum of at most $k$ simple fractions. It is a corollary of Theorem 14 below that this question must be answered negatively. Below, we will write $\mathcal{R}_0^d$ for the zero-totalized field of real numbers with the multiplicative inverse operation replaced by a division operation.

**Theorem 14.** For each $n \in \mathbb{N}^+$, let $p^n(x_1, \ldots, x_n) = 1 / x_1 + \ldots + 1 / x_n$. Then, for each $n \in \mathbb{N}^+$ with $n > 1$, there do not exist $n - 1$ simple fractions, say $q_1, \ldots, q_{n-1}$, such that $\mathcal{R}_0^d \models p^n(x_1, \ldots, x_n) = q_1 + \ldots + q_{n-1}$.

**Proof.** We prove this by induction on $n$.

The basis step, where $n = 2$, is easily proved by contradiction. Assume that there exists a simple fraction $q(x_1, x_2)$ such that $\mathcal{R}_0^d \models 1 / x_1 + 1 / x_2 = q(x_1, x_2)$. Then $\mathcal{R}_0^d \models 1 + 1 / x_2 = q(1, x_2)$. This contradicts the fact, following immediately from Theorem 9 that $\mathcal{Q}_0^d \not\models 1 + 1 / x_2 = q(1, x_2)$. Hence, there does not exist a simple fraction, say $q_1$, such that $\mathcal{R}_0^d \models p^2(x_1, x_2) = q_1$.

The inductive step is also proved by contradiction. Assume that there exist $n$ simple fractions, say $f_1 / g_1, \ldots, f_n / g_n$, such that $\mathcal{R}_0^d \models p^{n+1}(x_1, \ldots, x_{n+1}) = f_1 / g_1 + \ldots + f_n / g_n$. Here and in the remainder of the proof, all variables that have occurrences in $f_1, g_1, \ldots, f_n, g_n$ are understood to be among $x_1, \ldots, x_{n+1}$. Let $r_1, \ldots, r_n$ be real numbers, and let $r$ be the vector $(r_1, \ldots, r_n)$.

Let $p^{n+1}[r]$ be the unary function on real numbers that is the interpretation of $p^{n+1}(x_1, \ldots, x_{n+1})$ in $\mathcal{R}_0^d$ if $x_1, \ldots, x_n$ are assigned the values $r_1, \ldots, r_n$, respectively. Then $p^{n+1}[r](v) = u + 1 / v$ with $u = 1 / r_1 + \ldots + 1 / r_n$ a real number. Consequently, $p^{n+1}[r]$ is continuous everywhere except at 0.

For each $i \in \mathbb{N}^+$ with $i \leq n$, let $f_i[r], g_i[r]$, and $q_i[r]$ be the unary functions on real numbers that are the interpretations of $f_i, g_i$, and $f_i / g_i$, respectively, in $\mathcal{R}_0^d$ if $x_1, \ldots, x_n$ are assigned the values $r_1, \ldots, r_n$, respectively. Let $i \in \mathbb{N}^+$ be such
that $i \leq n$. Then $q_i[r](v) = f_i[r](v) / g_i[r](v)$. Because $g_i$ is a polynomial, $g_i[r]$ is continuous. Now assume that $g_i[r](0) \neq 0$. Then, in the case where $g_i[r](0) > 0$, by the continuity of $g_i[r]$, there exists a rational number $\epsilon_{i,r}$ such that, for each $v \in (-\epsilon_{i,r}, \epsilon_{i,r})$, $g_i[r](v) > 0$. Likewise, in the case where $g_i[r](0) < 0$, there exists a rational number $\epsilon_{i,r}$ such that, for each $v \in (-\epsilon_{i,r}, \epsilon_{i,r})$, $g_i[r](v) < 0$. From this, it follows that $g_i[r]$ is continuous in the interval $(-\epsilon_{i,r}, \epsilon_{i,r})$ if $g_i[r](0) \neq 0$.

From the continuity results established in the previous two paragraphs, we can prove by contradiction the claim that there exists an $i \in \mathbb{N}^+$ with $i \leq n$ such that $g_i[r](0) = 0$. Assume the contrary. For each $i \in \mathbb{N}^+$ with $i \leq n$, let $\epsilon_{i,r}$ be as indicated above. Moreover, let $\epsilon_r$ be the minimum of $\epsilon_{1,r}, \ldots, \epsilon_{n,r}$. Then, for each $i \in \mathbb{N}^+$ with $i \leq n$, $g_i[r]$ is continuous in the interval $(-\epsilon_r, \epsilon_r)$. From this and the fact that $p^{n+1}[r](v) = q_1[r](v) + \ldots + q_n[r](v)$ it follows that $p^{n+1}[r]$ is continuous in the interval $(-\epsilon_r, \epsilon_r)$. This contradicts the fact that $p^{n+1}[r]$ is not continuous at $0$. Hence, there exists an $i \in \mathbb{N}^+$ with $i \leq n$ such that $g_i[r](0) = 0$.

For each $i \in \mathbb{N}^+$ with $i \leq n$, let $f_i$ and $g_i$ be the terms obtained by replacing in $f_1$ and $g_1$, respectively, all occurrences of $x_{n+1}$ by $0$, and let $f_i[r]$ and $g_i[r]$ be the real numbers that are the interpretations of $f_i$ and $g_i$, respectively, in $\mathcal{R}_0$ if $x_1, \ldots, x_n$ are assigned the values $r_1, \ldots, r_n$, respectively. Then $\mathcal{R}_0 \models p^n(x_1, \ldots, x_n) = f_1[r] / g_1[r] + \ldots + f_n[r] / g_n[r]$ because $\mathcal{R}_0 \models 1 / 0 = 0$. Moreover, by the claim proved above, there exists an $i \in \mathbb{N}^+$ with $i \leq n$ such that $g_i[r](0) = 0$. From this, it follows that $g_i[r] \cdot \ldots \cdot g_n[r] = 0$.

Thus, we have established above that, for each vector $r$ of $n$ real numbers, $g_1[r] \cdot \ldots \cdot g_n[r] = 0$. Hence, $\mathcal{R}_0 \models g_1[r] \cdot \ldots \cdot g_n[r] = 0$. Using this result, we can prove the claim that there exists an $i \in \mathbb{N}^+$ with $i \leq n$ such that $\mathcal{R}_0 \models g_i = 0$. For each $i \in \mathbb{N}^+$ with $i \leq n$, $g_i$ is a multivariate polynomial in $n$ variables. Because the ring of polynomials over $\mathcal{R}$ in $n$ variables is an integral domain (see e.g. Proposition 4.29 in [14]), the ring of polynomials over $\mathcal{R}_0$ in $n$ variables is an integral domain as well. In other words, there are no zero divisors in this polynomial ring. From this and the fact that $\mathcal{R}_0 \models g_1 \cdot \ldots \cdot g_n = 0$, it follows that there exists an $i \in \mathbb{N}^+$ with $i \leq n$ such that $\mathcal{R}_0 \models g_i = 0$.

Take $j \in \mathbb{N}^+$ with $j \leq n$ such that $\mathcal{R}_0 \models g_j = 0$. Such a $j$ exists by the claim just proved. Then $\mathcal{R}_0 \models f_j[r] / g_j[r] = 0$. From this and the fact that $\mathcal{R}_0 \models p^n(x_1, \ldots, x_n) = f_1[r] / g_1[r] + \ldots + f_n[r] / g_n[r]$, it follows that $\mathcal{R}_0 \models p^n(x_1, \ldots, x_n) = f_1[r] / g_1[r] + \ldots + f_j[r] / g_j[r] + f_{j+1}[r] / g_{j+1}[r] + \ldots + f_n[r] / g_n[r]$. This contradicts the induction hypothesis. Hence, there do not exist $n$ simple fractions, say $f_1 / g_1, \ldots, f_n / g_n$, such that $\mathcal{R}_0 \models p^{n+1}(x_1, \ldots, x_{n+1}) = f_1 / g_1 + \ldots + f_n / g_n$. \hfill $\square$

Theorem [14] gives rise to the following corollary.

**Corollary 5.** For each $n \in \mathbb{N}^+$, let $p^n(x_1, \ldots, x_n) = 1 / x_1 + \ldots + 1 / x_n$. Then, for each $n \in \mathbb{N}^+$ with $n > 1$, there do not exist $n-1$ simple fractions, say $q_1, \ldots, q_{n-1}$, such that $E_{\mathcal{R}_0} \models p^n(x_1, \ldots, x_n) = q_1 + \ldots + q_{n-1}$.

**Proof.** This follows immediately from Theorem [14]. \hfill $\square$
Corollary 6. There does not exist a \( k \in \mathbb{N}^+ \) such that, for each term \( p \) over the signature \( \Sigma_{Md}^d \), there exists \( k \) simple fractions, say \( q_1, \ldots, q_k \), such that \( E_{Md}^d \vdash p = q_1 + \ldots + q_k \).

Proof. This follows immediately from Corollary 5. \( \square \)

6 Transformation of Closed Fractions

In this section, we establish results about the transformation into simple fractions for closed terms over the signature of divisive meadows.

The first result concerns the axioms of a divisive meadow.

Theorem 15. \( E_{Md}^d \) does not admit transformation into simple fractions for closed terms.

Proof. In order to prove this theorem by contradiction, assume that \( E_{Md}^d \) admits transformation into simple fractions for closed terms. Then, by Proposition 4, \( E_{Md}^d \vdash 1 + 1/2 = 1 \) or there exist \( n, m \in \mathbb{N}^+ \) such that \( E_{Md}^d \vdash 1 + 1/2 = n/m \) or \( E_{Md}^d \vdash 1 + 1/2 = -(n/m) \). However, \( E_{Md}^d \not\vdash 1 + 1/2 = 0 \) and, for all \( n, m \in \mathbb{N}^+ \), \( E_{Md}^d \not\vdash 1 + 1/2 = -(n/m) \) because \( Q_0 \not\vdash 1 + 1/2 = 0 \) and, for all \( n, m \in \mathbb{N}^+ \), \( Q_0 \not\vdash 1 + 1/2 = -(n/m) \). Consequently, there exist \( n, m \in \mathbb{N}^+ \) such that \( E_{Md}^d \vdash 1 + 1/2 = n/m \).

Let \( n, m \in \mathbb{N}^+ \) be such that \( E_{Md}^d \vdash 1 + 1/2 = n/m \). Then \( Q_0 \vdash 1 + 1/2 = n/m \).

Because also \( Q_0 \vdash 1 + 1/2 = 3/2 \) we have \( Q_0 \vdash n/m = 3/2 \). From this, it follows that \( Q_0 \vdash n \cdot 2 = m \cdot 3 \), and consequently, by Proposition 3, \( Q_0 \vdash n \cdot 2 = m \cdot 3 \).

Hence, \( n \cdot 2 = m \cdot 3 \). From this and the fact that \( 2 \) and \( 3 \) are relatively prime, \( n \) is a multiple of \( 3 \) and \( m \) is a multiple of \( 2 \).

Let \( m' \in \mathbb{Z} \) be such that \( m = 2 \cdot m' \). Then \( Md_2^d \vdash m = 2 \cdot m' \), and consequently, by Proposition 3, \( Md_2^d \vdash m = 2 \cdot m' \). From this and the fact that \( Md_2^d \vdash 2 = 0 \), it follows that \( Md_2^d \vdash m = 0 \). From this, it follows that \( Md_2^d \vdash n = m = 0 \). From this and the fact that \( Md_2^d \vdash 1 + 1/2 = 1 \), it follows that \( Md_2^d \vdash 1 + 1/2 = m = 0 \). This contradicts \( E_{Md}^d \vdash 1 + 1/2 = n/m \). Hence, \( E_{Md}^d \) does not admit transformation into simple fractions for closed terms. \( \square \)

The next result and the result following the second next result tell us that \( Q_0 \) is the only minimal model of \( E_{Md}^d \) with an infinite carrier that admits transformation into simple fractions for closed fractions. The second next result is an auxiliary result used to establish the result following it.

Theorem 16. \( Q_0^d \) admits transformation into simple fractions for closed terms.

Proof. Because of Theorem 5, it suffices to prove that for all \( p \in \mathcal{B} \), there exists a simple closed fraction \( q \) such that \( Q_0 \vdash p = q \). The proof is straightforward by induction on the structure of \( p \). If \( p \) is of the form \( \theta \), \( n/m \) or \( -(n/m) \), with \( n, m \in \mathbb{N}^+ \), then it is trivial to show that there exists a simple closed fraction \( q \) such that \( Q_0 \vdash p = q \). If \( p \) is of the form \( p' + q' \), then it follows immediately from the induction hypothesis and the claim that, for all \( n, m, n', m' \in \mathbb{N}^+ \), \( Q_0 \vdash n/m + n'/m' = (n \cdot m' + n' \cdot m) / (m \cdot m') \). This claim is easily proved using the fact that, for all \( n \in \mathbb{N}^+ \), \( Q_0 \vdash n/n = 1 \). \( \square \)
Lemma 1. Let $E \supseteq E_{Md}^{d}$ be a set of equations over the signature $\Sigma_{Md}^{d}$ and $p$ be a closed term over the signature $\Sigma_{Md}^{d}$. Then $E \not\vdash p = 0$ only if there exists a divisive meadow $M$ such that $M \models E \cup \{ p / p = 1 \}$ and $M \not\models 0 \neq 1$.

Proof. In order to prove this theorem by contraposition, assume that there does not exist a divisive meadow $M$ such that $M \models E \cup \{ p / p = 1 \}$ and $M \models 0 \neq 1$. In other words, assume that for each divisive meadow $M$, $M \models E \cup \{ p / p = 1 \}$ implies $M \models 0 = 1$.

Let $M'$ be a divisive cancellation meadow such that $M' \models p / p = 1$ implies $M' \models 0 = 1$. Assume that $M' \models p \not\neq 0$. Then $M' \models p / p = 1$ and, because $M' \models p / p = 1$ implies $M' \models 0 = 1$, also $M' \models 0 = 1$. However, if $M' \models p \not\neq 0$, then $M' \models p = 0$. This contradicts the assumption that $M' \models p \not\neq 0$. Hence, $M' \models p = 0$. From this, it follows that $E \vdash p = 0$ by Theorems 2 and 4. \hfill \Box

Theorem 17. An infinite minimal divisive meadow admits transformation into simple fractions for closed terms only if it is isomorphic to $Q_{0}^{d}$.

Proof. Let $M$ be an infinite minimal divisive meadow that admits transformation into simple fractions for closed terms. By Theorem 6, $Q_{0}^{d}$ is a homomorphic image of $M$. Let $E_M$ be the set of all equations of the form $u'/v' + u''/v'' = u/v$, with $u, v, u', v', u''$, and $v''$ of the form $\overline{n}$ or $-\overline{n}$ with $n \in \mathbb{N}^+$, that are satisfied by $M$.

In order to prove by contradiction that $M$ is isomorphic to $Q_{0}^{d}$, assume that $M$ is not isomorphic to $Q_{0}^{d}$. Then there exists a $l \in \mathbb{N}^+$ such that $M \not\models l/l = 1$ because $Q_{0}^{d}$ is the initial algebra of $E_{Md}^{d} \cup \{ l/l = 1 \mid l \in \mathbb{N}^+ \}$. Consequently, there exists a $l \in \mathbb{N}^+$ such that $E_{Md}^{d} \cup E_{M} \not\models l/l = 1$. Let $l \in \mathbb{N}^+$ be such that $E_{Md}^{d} \cup E_{M} \not\models l/l = 1$. By Lemma 1, there exists a divisive meadow $M'$ such that $M' \models E_{Md}^{d} \cup E_{M} \cup \{ (1 - l/l) / (1 - l/l) = 1 \}$ and $M' \models 0 \neq 1$.

Let $M'$ be a divisive meadow such that $M' \models E_{Md}^{d} \cup E_{M} \cup \{ (1 - l/l) / (1 - l/l) = 1 \}$ and $M' \models 0 \neq 1$, and let $M''$ be the minimal divisive submeadow of $M'$. Then, for each set of equations $E$ over $\Sigma_{Md}^{d}$, $M' \models E$ implies $M'' \models E$. In particular, $M'' \models E_{M}$. $M''$ does not have $Q_{0}^{d}$ as a homomorphic image because $Q_{0}^{d} \models (1 - l/l) / (1 - l/l) = 0$. From this and Theorem 6, it follows that $M''$ is a finite minimal divisive meadow. We know from Lemmas 4.1 and 4.8 in [5] and Theorem 4 in this paper that each finite minimal divisive meadow is isomorphic to a divisive meadow $Md_{k}^{d}$ for some $k \in \mathbb{N}^+$ that is square-free. So we may assume that $M'' = Md_{k}^{d}$ for some $k \in \mathbb{N}^+$.

Let $k \in \mathbb{N}^+$ be such that $M'' = Md_{k}^{d}$. Then there exist $n, m \in \mathbb{Z}$ such that $1 + 1/k = \overline{n}/m \in E_{M}$. Let $n, m \in \mathbb{Z}$ be such that $1 + 1/k = \overline{n}/m \in E_{M}$. Then $Q_{0}^{d} \models 1 + 1/k = \overline{n}/m$ because $Q_{0}^{d}$ is a homomorphic image of $M$. From this, it follows that $Q_{0}^{d} \models (k + 1)/k = n/m$. Consequently, by Proposition 3, $Q_{0}^{d} \models k + 1/k = \overline{n}/m$. From this, it follows that $Q_{0}^{d} \models n \cdot k = \overline{m} \cdot (k + 1)$ and consequently, by Proposition 3, $Q_{0}^{d} \models n \cdot k = m \cdot (k + 1)$. Hence, $n \cdot k = m \cdot (k + 1)$. From this and the fact that $k$ and $k + 1$ are relatively prime, it follows that $n$ is a multiple of $k + 1$ and $m$ is a multiple of $k$. Let $m' \in \mathbb{Z}$ be such that $m = k \cdot m'$. Then $Md_{k}^{d} \models m = k \cdot m'$, and consequently, by Proposition 3.
\( \mathcal{M}d_k^d \models m = k \cdot m'. \) From this and the fact that \( \mathcal{M}d_k^d \models k = 0, \) it follows that \( \mathcal{M}d_k^d \models m = 0. \) From this, it follows that \( \mathcal{M}d_k^d \models n / m = 0. \) From this and the fact that \( \mathcal{M}d_k^d \models 1 + 1 / k = 1, \) it follows that \( \mathcal{M}d_k^d \not\models 1 + 1 / k = n / m. \) Because \( \mathcal{M}d_k^d = \mathcal{M}'' \) and \( 1 + 1 / k = n / m \in E_{\mathcal{M}}, \) this contradicts \( \mathcal{M}'' \models E_{\mathcal{M}}. \) Hence, \( \mathcal{M} \) is isomorphic to \( Q_0^d. \)

The next theorem tells us that admitting transformation into simple fractions for closed terms is a property of divisive meadows that cannot be expressed as a first-order theory.

**Theorem 18.** Admitting transformation into simple fractions for closed terms is not an elementary property of divisive meadows.

**Proof.** In order to prove this theorem by contradiction, assume that there exists a first-order theory \( T \) over \( \Sigma_{\mathcal{M}d}^d \) such that the models of \( E_{\mathcal{M}d}^d \cup T \) are precisely the divisive meadows with the mentioned property. In this proof, we write \( NC \) for \( \exists x \cdot (x \neq 0 \land x / x \neq 1). \)

Let \( T \) be a first-order theory over \( \Sigma_{\mathcal{M}d}^d \) such that the models of \( E_{\mathcal{M}d}^d \cup T \) are precisely the divisive meadows with the mentioned property. Because of Theorem 7, for each \( k' \in \mathbb{N}, \) there exists a square-free \( k'' \in \mathbb{N} \) with \( k'' \) not prime and \( k'' > k' \) such that \( \mathcal{M}d_{k''}^d \models E_{\mathcal{M}d}^d \cup T \cup \{ \text{NC} \} \cup \{ k \neq 0 \mid k \in \mathbb{N}^+, k \leq k' \}. \) From this, it follows that, for each \( k' \in \mathbb{N}, \) there exists a divisive meadow \( \mathcal{M} \) such that \( \mathcal{M} \models E_{\mathcal{M}d}^d \cup T \cup \{ \text{NC} \} \cup \{ k \neq 0 \mid k \in \mathbb{N}^+, k \leq k' \}. \) Hence, by the compactness of first-order logic, there exists a divisive meadow \( \mathcal{M}' \) such that \( \mathcal{M}' \models E_{\mathcal{M}d}^d \cup T \cup \{ \text{NC} \} \cup \{ k \neq 0 \mid k \in \mathbb{N}^+, k < k' \}. \) In other words, there exists a divisive meadow \( \mathcal{M} \) of characteristic 0 such that \( \mathcal{M} \models E_{\mathcal{M}d}^d \cup T \cup \{ \text{NC} \}. \) This contradicts Theorem 17 because a divisive meadow of characteristic 0 is infinite and a divisive meadow that satisfies \( NC \) is not a divisive cancellation meadow. Hence, there does not exist a first-order theory \( T \) over \( \Sigma_{\mathcal{M}d}^d \) such that the models of \( E_{\mathcal{M}d}^d \cup T \) are exactly the divisive meadows with the mentioned property. \( \square \)

**7 Miscellaneous Results about Divisive Meadows**

In this section, we establish two results that are related to the results in preceding sections, but do not concern fractions.

In the proof of Theorem 17, one of the main results of this paper, Theorem 6 plays an important part. The first result, which is a generalization of the result that a polynomial in canonical form is derivably equal to 0 only if each of its coefficients is derivably equal to 0, is established by means of Theorem 6 as well.

**Theorem 19.** Let \( f(x) \) be a term over \( \Sigma_{\mathcal{M}d}^d \) of the form \( a_n \cdot x^n + \ldots + a_1 \cdot x + a_0, \) where \( a_i \) is a closed term over \( \Sigma_{\mathcal{M}d}^d \) for each \( i \in \{0, \ldots, n\}. \) Then \( E_{\mathcal{M}d}^d \models f(x) = 0 \) only if, for each \( i \in \{0, \ldots, n\}, \) \( E_{\mathcal{M}d}^d \models a_i = 0. \)

**Proof.** Because \( E_{\mathcal{M}d}^d \models f(x) = 0, \) we have \( Q_0^d \models f(x) = 0. \) From this and the fact that there exists an \( i \in \{0, \ldots, n\} \) such that \( Q_0^d \not\models a_i = 0 \) only if \( Q_0^d \not\models f(x) = 0, \) it follows that \( Q_0^d \models a_i = 0 \) for each \( i \in \{0, \ldots, n\}. \)
In order to prove this theorem by contradiction, assume that there exists an
\( i \in \{0, \ldots, n\} \) such that \( E_{\text{Md}}^d \Downarrow a_i = 0 \). Substitution of \( x \) by 0 in \( f(x) = 0 \) yields
\( E_{\text{Md}}^d \Downarrow a_0 = 0 \). Therefore, there exists an \( i \in \{0, \ldots, n\} \) such that \( E_{\text{Md}}^d \Downarrow a_i = 0 \). Let
\( m \in \{1, \ldots, n\} \) be the maximal \( i \in \{1, \ldots, n\} \) such that \( E_{\text{Md}}^d \Downarrow a_i = 0 \). Then
\( E_{\text{Md}}^d \Downarrow a_m \cdot x^m + \ldots + a_1 \cdot x = 0 \).

By Lemma 1 we know that there exists a divisive meadow \( \mathcal{M} \) such that
\( \mathcal{M} \models E_{\text{Md}}^d \Downarrow \{a_m / a_m = 1\} \) and \( \mathcal{M} \models 0 \neq 1 \). Let \( \mathcal{M} \) be a divisive meadow such that
\( \mathcal{M} \models E_{\text{Md}}^d \Downarrow \{a_m / a_m = 1\} \) and \( \mathcal{M} \models 0 \neq 1 \), and let \( \mathcal{M}' \) be the minimal divisive submeadow of \( \mathcal{M} \). Then, for each equation \( \phi \) over \( \Sigma_{\text{Md}}, \mathcal{M} \models \phi \) implies
\( \mathcal{M}' \models \phi \). In particular, \( \mathcal{M}' \models a_m / a_m = 1 \).

\( \mathcal{Q}_0^d \) is not a homomorphic image of \( \mathcal{M}' \) because otherwise \( \mathcal{Q}_0^d \models a_m / a_m = 1 \) and consequently \( \mathcal{Q}_0^d \models a_m \neq 0 \). From this and Theorem 20 it follows that \( \mathcal{M}' \) is
a finite minimal divisive meadow. We know from Lemmas 4.1 and 4.8 in [5] and
Theorem 4 in this paper that each finite minimal divisive meadow is isomorphic
to a divisive submeadow of \( \mathcal{M} \).

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Theorem 4 in this paper that each finite minimal divisive meadow is isomorphic
to a divisive submeadow of \( \mathcal{M} \).

A polynomial in canonical form is derivably equal to 0 only if each of its closed
substitution instances is derivably equal to 0. The question arises whether “only
if” can be replaced by “if and only if”. The proof of the next theorem gives a
negative answer this question.

**Theorem 20.** \( E_{\text{Md}}^d \) is not \( \omega \)-complete.

**Proof.** First we prove by contradiction that \( E_{\text{Md}}^d \Downarrow (1 - \frac{2}{2}) \cdot (x^2 - x) = 0 \) and next we prove that
\( E_{\text{Md}}^d \Downarrow (1 - \frac{2}{2}) \cdot (p^2 - p) = 0 \) for each closed term \( p \) over \( \Sigma_{\text{Md}} \).

Assume that \( E_{\text{Md}}^d \Downarrow (1 - \frac{2}{2}) \cdot (p^2 - p) = 0 \), and consequently \( \mathcal{M} \models x^2 = x \) for each divisive meadow \( \mathcal{M} \) of characteristic 2.

Let \( \mathcal{M}' \) be a finite algebraic extension of \( \mathcal{M}_d \) such that \( \mathcal{M}' \models x^2 + x + 1 = 0 \). Then \( \mathcal{M}' \) is a divisive meadow of characteristic 2 such that \( \mathcal{M}' \not\models x^2 = x \). This is easy to see: otherwise \( \mathcal{M}' \models x + x + 1 = 0 \) and consequently \( \mathcal{M}' \models 1 = 0 \). However, \( \mathcal{M}' \not\models \frac{x^2}{x} = x \) contradicts \( \mathcal{M} \models \frac{x^2}{x} = x \) for each divisive meadow \( \mathcal{M} \) of characteristic 2. Hence, \( E_{\text{Md}}^d \Downarrow (1 - \frac{2}{2}) \cdot (x^2 - x) = 0 \).
Let $M$ be a divisive cancellation meadow. If $M \models p = 0$ or $M \models p = 1$, then it is easily proved by structural induction that either $M \models p^2 - p = 0$ or $M \models p^2 - p = 1$, and consequently $M \models p^2 - p = 0$, for each closed term $p$ over $\Sigma_{Md}^d$. If $M \models p \neq 0$, then $M \models 1 - \frac{2}{2} = 0$. Hence, $M \models (1 - \frac{2}{2}) \cdot (p^2 - p) = 0$ for each closed term $p$ over $\Sigma_{Md}^d$. From this, it follows that $E_{Md}^d \vdash (1 - \frac{2}{2}) \cdot (p^2 - p) = 0$ for each closed term $p$ over $\Sigma_{Md}^d$ by Theorems $2$ and $4$.

The $\omega$-completeness question can also be posed about extensions of $E_{Md}^d$ that exclude divisive meadows of non-zero characteristic, such as $E_{Md}^d \cup \{(1 + x^2 + y^2) / (1 + x^2 + y^2) = 1 \}$ ($Q_0^d$ is the initial algebra among the divisive meadows that satisfy this extension of $E_{Md}^d$). Such variants of the question are related to Hilbert’s tenth problem and are harder to answer.

8 Concluding Remarks

We have shown that the setting of meadows allows workable syntactic definitions of a fraction and a simple fraction to be given. This only means that we have a point of departure for the development of a workable theory about fractions. We have made a start with the development of such a theory, but there remain many open questions. For instance, it is an open question, arising from Theorem $11$ whether each divisive meadow of non-zero characteristic for which there exists an $n \in \mathbb{N}$ such that each element of its carrier is the root of a non-trivial polynomial of degree $n$ or less admits transformation into simple fractions — we know already from Theorem $11$ that there exists such an $n$ for each divisive meadow of prime characteristic. Another open question, arising from the proof of Corollary $6$ is whether there exists a natural number $k$ such that each term in one variable over $\Sigma_{Md}^d$ is derivably equal to a sum of at most $k$ simple fractions.

There are questions that are not complicated for simple fractions, but complicated for more complex terms. Let $C_0^d$ be the zero-totalized field of complex numbers with the multiplicative inverse operation replaced by a division operation. Using Robinson’s classical result that the first order theory of an algebraically closed field is model complete [20], it is proved in [3] that the equational theory of $C_0^d$ and the equational theory of the class of all models of $E_{Md}^d \cup \{a \cdot \frac{n}{n} = 1 \mid n \in \mathbb{N}^+ \}$ are the same. From this, it follows that $C_0^d \models p = 0$ if and only if $E_{Md}^d \cup \{a \cdot \frac{n}{n} = 1 \mid n \in \mathbb{N}^+ \} \vdash p = 0$. A simple direct proof with the theory developed so far can be found if $p$ is restricted to simple fractions. However, it seems less straightforward to find such a proof if $p$ is restricted to sums of simple fractions.

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