Graph Treewidth and Geometric Thickness Parameters*

Vida Dujmović† David R. Wood‡

March 30, 2005; Revised: September 25, 2018

Abstract

Consider a drawing of a graph $G$ in the plane such that crossing edges are coloured differently. The minimum number of colours, taken over all drawings of $G$, is the classical graph parameter thickness. By restricting the edges to be straight, we obtain the geometric thickness. By additionally restricting the vertices to be in convex position, we obtain the book thickness. This paper studies the relationship between these parameters and treewidth.

Our first main result states that for graphs of treewidth $k$, the maximum thickness and the maximum geometric thickness both equal $\lceil k/2 \rceil$. This says that the lower bound for thickness can be matched by an upper bound, even in the more restrictive geometric setting. Our second main result states that for graphs of treewidth $k$, the maximum book thickness equals $k$ if $k \leq 2$ and equals $k + 1$ if $k \geq 3$. This refutes a conjecture of Ganley and Heath [Discrete Appl. Math. 109(3):215–221, 2001]. Analogous results are proved for outerthickness, arboricity, and star-arboricity.

---

1This research was initiated in 2005 while both authors were in the School of Computer Science at McGill University, Montréal, Canada. A preliminary version of this paper was published in the Proceedings of the 13th International Symposium on Graph Drawing (GD ’05), Lecture Notes in Computer Science 3843:129–140, Springer, 2006. The full version was published in Discrete and Computational Geometry 37.4:641–670, 2007. That version contained a false conjecture, which is corrected on page 26 of this version. No other changes have been made.

2Department of Mathematics and Statistics, McGill University, Montréal, Canada. Partially supported by NSERC, Centre de Recherches Mathématiques (CRM), and Institut des Sciences Mathématiques (ISM). Now at: School of Computer Science and Electrical Engineering, University of Ottawa, Ottawa, Canada (vida.dujmovic@uottawa.ca). Research supported by NSERC.

3Departament de Matemàtica Aplicada II, Universitat Politècnica de Catalunya, Barcelona, Spain. Supported by a Marie Curie Fellowship of the European Community under contract 023865, and by the projects MCYT-FEDER BFM2003-00368 and Gen. Cat 2001SGR00224. Now at: School of Mathematical Sciences, Monash University, Melbourne, Australia (david.wood@monash.edu). Research supported by the Australian Research Council.
1 Introduction

Partitions of the edge set of a graph into a small number of ‘nice’ subgraphs are in the mainstream of graph theory. For example, in a proper edge colouring, the subgraphs of the partition are matchings. If each subgraph of a partition is required to be planar (respectively, outerplanar, a forest, a star-forest), then the minimum number of subgraphs in a partition of a graph $G$ is the thickness (outerthickness, arboricity, star-arboricity) of $G$. Thickness and arboricity are classical graph parameters that have been studied since the early 1960s.

The first results in this paper concern the relationship between the above parameters and treewidth, which is a more modern graph parameter that is particularly important in structural and algorithmic graph theory; see the surveys [16, 66]. In particular, we determine the maximum thickness, maximum outerthickness, maximum arboricity, and maximum star-arboricity of a graph with treewidth $k$. These results are presented in Section 3 (following some background graph theory in Section 2).

The main results of the paper are about graph partitions with an additional geometric property. Namely, there is a drawing of the graph, and each subgraph in the partition is drawn without crossings. This type of drawing has applications in graph visualisation (where each noncrossing subgraph is coloured by a distinct colour), and in multilayer VLSI (where each noncrossing subgraph corresponds to a set of wires that can be routed without crossings in a single layer). With no restriction on how the edges are drawn, the minimum number of noncrossing subgraphs, taken over all drawings of $G$, is again the thickness of $G$. By restricting the edges to be drawn straight, we obtain the geometric thickness of $G$. By further restricting the vertices to be in convex position, we obtain the book thickness of $G$. These geometric parameters are introduced in Section 4.

Our main results determine the maximum geometric thickness and maximum book thickness of a graph with treewidth $k$. Analogous results are proved for geometric variations of outerthickness, arboricity, and star-arboricity. These geometric results are stated in Section 5. The general approach that is used in the proofs of our geometric upper bounds is described in Section 6. The proofs of our geometric results are in Sections 7–9. Section 10 concludes with numerous open problems.

2 Background Graph Theory

For undefined graph-theoretic terminology, see the monograph by Diestel [25]. We consider graphs $G$ that are simple, finite, and undirected. Let $V(G)$ and $E(G)$ respectively denote the vertex and edge sets of $G$. For $A, B \subseteq V(G)$, let $G[A; B]$ denote the bipartite subgraph of $G$ with vertex set $A \cup B$ and edge set $\{vw \in E(G) : v \in A, w \in B\}$.

A partition of a graph $G$ is a proper partition $\{E_1, E_2, \ldots, E_t\}$ of $E(G)$; that is, $\bigcup\{E_i : 1 \leq i \leq t\} = E(G)$ and $E_i \cap E_j = \emptyset$ whenever $i \neq j$. Each part $E_i$ can be thought of as a spanning subgraph $G_i$ of $G$ with $V(G_i) := V(G)$ and $E(G_i) := E_i$. We also consider a partition to be an edge-colouring, where each edge in $E_i$ is coloured $i$. In an edge-coloured graph, a vertex $v$ is colourful if all the edges incident to $v$ receive distinct colours.

A graph parameter is a function $f$ such that $f(G) \in \mathbb{N} := \{0, 1, 2, \ldots\}$ for every graph.
Let $\mathcal{G}$ be a graph class. For a graph $G$, let $f(G) := \max\{f(G) : G \in \mathcal{G}\}$. If $f(G)$ is unbounded, we write $f(G) := \infty$.

Our interest is in drawings of graphs in the plane; see [19, 23, 51, 62, 69]. A drawing $\phi$ of graph $G$ is a pair $(\phi_V, \phi_E)$, where:

- $\phi_V$ is an injection from the vertex set $V(G)$ into $\mathbb{R}^2$, and
- $\phi_E$ is a mapping from the edge set $E(G)$ into the set of simple curves in $\mathbb{R}^2$, such that for each edge $vw \in E(G)$,
  - the endpoints of the curve $\phi_E(vw)$ are $\phi_V(v)$ and $\phi_V(w)$, and
  - $\phi_V(x) \not\in \phi_E(vw)$ for every vertex $x \in V(G) \setminus \{v, w\}$.

If $H$ is a subgraph of a graph $G$, then every drawing $\phi$ of $G$ induces a subdrawing of $H$ obtained by restricting the functions $\phi_V$ and $\phi_E$ to the elements of $H$. Where there is no confusion, we do not distinguish between a graph element and its image in a drawing.

A set of points $S \subset \mathbb{R}^2$ is in general position if no three points in $S$ are collinear. A general position drawing is one in which the vertices are in general position.

Two edges in a drawing cross if they intersect at some point other than a common endpoint. A cell of a drawing $\phi$ of $G$ is a connected component of $\mathbb{R}^2 \setminus \{\phi_V(v) : v \in V(G)\} \cup \{\phi_E(vw) : vw \in E(G)\}$. Thus each cell of a drawing is an open subset of $\mathbb{R}^2$ bounded by edges, vertices, and crossing points. Observe that a drawing of a (finite) graph has exactly one cell of infinite measure, called the outer cell. A graph drawing with no crossings is noncrossing. A graph that admits a noncrossing drawing is planar. A drawing in which all the vertices are on the boundary of the outer cell is outer. A graph that admits an outer noncrossing drawing is outerplanar.

The thickness of a graph $G$, denoted by $\theta(G)$, is the minimum number of planar subgraphs that partition $G$. Thickness was first defined by Tutte [73]; see the surveys [46, 60]. The outerthickness of a graph $G$, denoted by $\theta_o(G)$, is the minimum number of outerplanar subgraphs that partition $G$. Outerthickness was first studied by Guy [40]; also see [31, 33, 41, 42, 52, 65].

The arboricity of a graph $G$, denoted by $a(G)$, is the minimum number of forests that partition $G$. Nash-Williams [61] proved that
\[
a(G) = \max_{H \subseteq G} \left[ \frac{|E(H)|}{|V(H)| - 1} \right].
\]

A star is a tree with diameter at most 2. A star-forest is a graph in which each component is a star. The star-arboricity of a graph $G$, denoted by $sa(G)$, is the minimum number of star-forests that partition $G$. Star arboricity was first studied by Akiyama and Kano [1]; also see [3, 5, 39, 43, 47].

It is well known that thickness, outerthickness, arboricity, and star-arboricity are within a constant factor of each other. In particular, Gonçalves [38] recently proved a longstanding

---

1 A simple curve is a homeomorphic image of the closed unit interval; see [59] for background topology.

2 In the literature on crossing numbers it is customary to require that intersecting edges cross ‘properly’ and do not ‘touch’. This distinction will not be important in this paper.
conjecture that every planar graph $G$ has outerthickness $\theta_o(G) \leq 2$. Thus $\theta_o(G) \leq 2 \cdot \theta(G)$ for every graph $G$. Every planar graph $G$ satisfies $|E(G)| < 3(|V(G)| - 1)$. Thus $a(G) \leq 3 \cdot \theta(G)$ for every graph $G$ by Equation (1). Similarly, every outerplanar graph $G$ satisfies $|E(G)| \leq 2(|V(G)| - 1)$. Thus $a(G) \leq 2 \cdot \theta_o(G)$ for every graph $G$ by Equation (1). Hakimi et al. [43] proved that every outerplanar graph $G$ has star-arboricity $sa(G) \leq 3$, and that every planar graph $G$ has star-arboricity $sa(G) \leq 5$. (Algor and Alon [3] constructed planar graphs for which $sa(G) = 5$.) Thus $sa(G) \leq 3 \cdot \theta_o(G)$ and $sa(G) \leq 5 \cdot \theta(G)$ for every graph $G$. It is easily seen that every tree $G$ has star-arboricity $sa(G) \leq 2$. Thus $sa(G) \leq 2 \cdot a(G)$ for every graph $G$. Summarising, we have the following set of inequalities.

$$\theta(G) \leq \theta_o(G) \leq a(G) \leq sa(G) \leq 5 \cdot \theta(G) .$$

$$2 \cdot \theta_o(G) \leq 2 \cdot a(G) \leq 3 \cdot \theta(G) \leq 3 \cdot \theta_o(G) \leq 5 \cdot \theta(G) .$$

Let $K_n$ be the complete graph on $n$ vertices. A set of $k$ pairwise adjacent vertices in a graph $G$ is a $k$-clique. For a vertex $v$ of $G$, let $N_G(v) := \{w \in V(G) : vw \in E(G)\}$ and $N_G[v] := N_G(v) \cup \{v\}$. We say $v$ is $k$-simplicial if $N_G(v)$ is a $k$-clique (and hence $N_G[v]$ is a $(k + 1)$-clique).

For each integer $k \geq 1$, a $k$-tree is a graph $G$ such that either:

- $G \cong K_{k+1}$, or
- $G$ has a $k$-simplicial vertex $v$ and $G \setminus v$ is a $k$-tree.

Suppose that $C$ is a clique in a graph $G$, and $S$ is a nonempty set with $S \cap V(G) = \emptyset$. Let $G'$ be the graph with vertex set $V(G') := V(G) \cup S$, and edge set $E(G') := E(G) \cup \{vx : v \in S, x \in C\}$. We say that $G'$ is obtained from $G$ by adding $S$ onto $C$. If $S = \{v\}$ then $G'$ is obtained from $G$ by adding $v$ onto $C$. Observe that if $|C| = k$, and $G$ is a $k$-tree or $G \cong K_k$, then $G'$ is a $k$-tree.

The treewidth of a graph $G$ is the minimum $k \in \mathbb{N}$ such that $G$ is a spanning subgraph of a $k$-tree. Let $\mathcal{T}_k$ be the class of graphs with treewidth at most $k$. Many families of graphs have bounded treewidth; see [10]. $\mathcal{T}_1$ is the class of forests. Graphs in $\mathcal{T}_2$ are obviously planar—a 2-simplicial vertex can always be drawn near the edge connecting its two neighbours without introducing a crossing. Graphs in $\mathcal{T}_2$ are characterised as those with no $K_4$-minor, and are sometimes called series-parallel.

### 3 Abstract Parameters and Treewidth

In this section we determine the maximum value of each of thickness, outerthickness, arboricity, and star-arboricity for graphs of treewidth $k$. Since every graph with treewidth $k$ is a subgraph of a $k$-tree, to prove the upper bounds we need only consider $k$-trees. The proofs of the lower bounds employ the complete split graph $K_{k,s}^*$ (for $k, s \geq 1$), which is the $k$-tree obtained by adding a set $S$ of $s$ vertices onto an initial $k$-clique $K$; see Figure [1]

Suppose that the edges of $K_{k,s}^*$ are coloured $1, 2, \ldots, \ell$. Let $c(e)$ be the colour assigned to each edge $e$ of $K_{k,s}^*$. The colour vector of each vertex $v \in S$ is the set $\{(c(uv), u) : u \in K\}$. Note that there are $\ell^k$ possible colour vectors.
Proposition 1. The maximum thickness of a graph in $T_k$ is $\lceil k/2 \rceil$; that is,

$$\theta(T_k) = \lceil k/2 \rceil.$$  

Proof. First we prove the upper bound. Ding et al. [27] proved that for all $k_1, k_2, \ldots, k_t \in \mathbb{N}$ with $k_1 + k_2 + \cdots + k_t = k$, every $k$-tree $G$ can be partitioned into $t$ subgraphs $G_1, G_2, \ldots, G_t$, such that each $G_i$ is a $k_i$-tree. Note that the $t = 2$ case, which implies the general result, was independently proved by Chhajed [20]. With $k_i = 2$, and since 2-trees are planar, we have $\theta(G) \leq \lceil k/2 \rceil$. (Theorem 1 provides an alternative proof with additional geometric properties.)

Now we prove the lower bound. If $k \leq 2$ then $\theta(T_k) \geq \theta(K_2) = 1 = \lceil k/2 \rceil$. Now assume that $k \geq 3$. Let $\ell := \lceil k/2 \rceil - 1$ and $s := 2\ell^k + 1$. Thus $\ell \geq 1$. Suppose that $\theta(K_{k,s}^\ast) \leq \ell$. In the corresponding edge $\ell$-colouring of $K_{k,s}^\ast$, there are $\ell^k$ possible colour vectors. Thus there are at least three vertices $x, y, z \in S$ with the same colour vector. At least $\lceil \ell/k \rceil \geq 2$ of the $k$ edges incident to $x$ are assigned the same colour. Say these edges are $xa, xb, xc$. Since $y$ and $z$ have the same colour vector as $x$, the $K_{3,3}$ subgraph induced by $\{xa, xb, xc, ya, yb, yc, za, zb, zc\}$ is monochromatic. This is a contradiction since $K_{3,3}$ is not planar. Thus $\theta(T_k) \geq \theta(K_{k,s}^\ast) \geq \ell + 1 = \lceil k/2 \rceil$. \hfill $\Box$

Proposition 2. The maximum arboricity of a graph in $T_k$ is $k$; that is,

$$a(T_k) = k.$$  

Proof. By construction, $|E(G)| = k|V(G)| - k(k + 1)/2$ for every $k$-tree $G$. It follows from Equation (1) that $a(G) \leq k$, and $a(G) = k$ if $|V(G)|$ is large enough. \hfill $\Box$

Proposition 3. The maximum outerthickness of a graph in $T_k$ is $k$; that is,

$$\theta_o(T_k) = k.$$  

Proof. Since a forest is outerplanar, $\theta_o(T_k) \leq a(T_k) = k$ by Proposition 2. Now we prove the lower bound. If $k = 1$ then $\theta_o(T_k) \geq \theta_o(K_2) = 1$. Now assume that $k \geq 2$. Let $\ell := k - 1$ and $s := 2\ell^k + 1$. Then $\ell \geq 1$. Suppose that $\theta_o(K_{k,s}^\ast) \leq \ell$. In the corresponding edge $\ell$-colouring of $K_{k,s}^\ast$, there are $\ell^k$ possible colour vectors. Thus there are at least three vertices $x, y, z \in S$ with the same colour vector. At least $\lceil k/\ell \rceil = 2$ of the $k$ edges incident to $x$ are assigned the
same colour. Say these edges are \( xa \) and \( xb \). Since \( y \) and \( z \) have the same colour vector as \( x \), the \( K_{2,3} \) subgraph induced by \{\( xa, xb, ya, yb, za, zb \)\} is monochromatic. This is a contradiction since \( K_{2,3} \) is not outerplanar. Thus \( \theta_0(T_k) \geq \theta_0(K^*_{k,s}) \geq \ell + 1 = k. \)

**Proposition 4.** The maximum star-arboricity of a graph in \( T_k \) is \( k + 1 \); that is,

\[
\text{sa}(T_k) = k + 1.
\]

**Proof.** The upper bound \( \text{sa}(T_k) \leq k + 1 \) was proved by Ding et al. \[27\]. For the lower bound, let \( s := k^k + 1 \). Let \( G \) be the graph obtained from the \( k \)-tree \( K^*_{k,s} \) by adding, for each vertex \( v \in S \), one new vertex \( v' \) onto \{\( v \)\}. Clearly \( G \) has treewidth \( k \). Suppose that \( \text{sa}(G) \leq k \). In the corresponding edge \( k \)-colouring of \( K^*_{k,s} \) there are \( k^k \) possible colour vectors. Since \( |S| > k^k \), there are two vertices \( x, y \in S \) with the same colour vector. No two edges in \( G[\{x\}; K] \) receive the same colour, as otherwise, along with \( y \), we would have a monochromatic 4-cycle. Thus all \( k \) colours are present on the edges of \( G[\{x\}; K] \) and \( G[\{y\}; K] \). Let \( p \) be the vertex in \( K \) such that \( xp \) and \( yp \) receive the same colour as \( xx' \). Thus \((x', x, p, y)\) is a monochromatic 4-vertex path, which is not a star. This contradiction proves that \( \text{sa}(T_k) \geq \text{sa}(G) \geq k + 1. \)

## 4 Geometric Parameters

The *thickness* of a graph drawing is the minimum \( k \in \mathbb{N} \) such that the edges of the drawing can be partitioned into \( k \) noncrossing subdrawings; that is, each edge is assigned one of \( k \) colours such that edges with same colour do not cross. Every planar graph can be drawn with its vertices at prespecified locations \[44, 64\]. Thus a graph with thickness \( k \) has a drawing with thickness \( k \) \[44\]. However, in such a drawing the edges might be highly curved. This motivates the notion of geometric thickness.

A drawing \((\phi_V, \phi_E)\) of a graph \( G \) is *geometric* if the image of each edge \( \phi_E(vw) \) is a straight line-segment (by definition, with endpoints \( \phi_V(v) \) and \( \phi_V(w) \)). Thus a geometric drawing of a graph is determined by the positions of its vertices. We thus refer to \( \phi_V \) as a geometric drawing.

The *geometric thickness* of a graph \( G \), denoted by \( \overline{\ell}(G) \), is the minimum \( k \in \mathbb{N} \) such that there is a geometric drawing of \( G \) with thickness \( k \). Kainen \[50\] first defined geometric thickness under the name of *real linear thickness*, and it has also been called *rectilinear thickness*. By the Fáry-Wagner theorem \[35, 74\], a graph has geometric thickness 1 if and only if it is planar. Graphs of geometric thickness \( 2 \), the so-called *doubly linear* graphs, were studied by Hutchinson et al. \[48\].

The *outerthickness* (respectively, *arboricity, star-arboricity*) of a graph drawing is the minimum \( k \in \mathbb{N} \) such that the edges of the drawing can be partitioned into \( k \) outer noncrossing

---

\[\text{Lemma 2}\] provides an alternative proof that \( \text{sa}(T_k) \leq k + 1 \). The same result can be concluded from a result by Hakimi et al. \[43\]. A vertex colouring with no bichromatic edge and no bichromatic cycle is *acyclic*. It is folklore that every \( k \)-tree \( G \) has an acyclic \((k + 1)\)-colouring \[36\]. (Proof. If \( G \approx K_{k+1} \) then the result is trivial. Otherwise, let \( v \) be a \( k \)-simplicial vertex. By induction, \( G \setminus v \) has an acyclic \((k + 1)\)-colouring. One colour is not present on the \( k \) neighbours of \( v \). Give this colour to \( v \). Thus there is no bichromatic edge. The neighbours of \( v \) have distinct colours since they form a clique. Thus there is no bichromatic cycle.) Hakimi et al. \[43\] proved that a graph with an acyclic \( c \)-colouring has star arboricity at most \( c \). Thus \( \text{sa}(T_k) \leq k + 1. \)
subdrawings (noncrossing forests, noncrossing star-forests). Again a graph with outerthickness (arboricity, star-arboricity) $k$ has a drawing with outerthickness (arboricity, star-arboricity) $k$ \cite{44, 64}. We generalise the notion of geometric thickness as follows. The geometric outerthickness (geometric arboricity, geometric star-arboricity) of a graph $G$, denoted by $\theta_o(G)$ ($\overline{\theta}(G)$, $\overline{sa}(G)$), is the minimum $k \in \mathbb{N}$ such that there is a geometric drawing of $G$ with outerthickness (arboricity, star-arboricity) $k$.

A geometric drawing in which the vertices are in convex position is called a book embedding. The book thickness of a graph $G$, denoted by $bt(G)$, is the minimum $k \in \mathbb{N}$ such that there is book embedding of $G$ with thickness $k$. Note that whether two edges cross in a book embedding is simply determined by the relative positions of their endpoints in the cyclic order of the vertices around the convex hull. A book embedding with thickness $k$ is commonly called a $k$-page book embedding: one can think of the vertices as being ordered on the spine of a book and each noncrossing subgraph being drawn without crossings on a single page. Book embeddings, first defined by Ollmann \cite{63}, are ubiquitous structures with a variety of applications; see \cite{28} for a survey with over 50 references. A book embedding is also called a stack layout, and book thickness is also called stacknumber, pagenumber and fixed outerthickness.

A graph has book thickness 1 if and only if it is outerplanar \cite{13}. Bernhart and Kainen \cite{13} proved that a graph has a book thickness at most 2 if and only if it is a subgraph of a Hamiltonian planar graph. Yannakakis \cite{78} proved that every planar graph has book thickness at most 4.

The book arboricity (respectively, book star-arboricity) of a graph $G$, denoted by $ba(G)$ ($bsa(G)$), is the minimum $k \in \mathbb{N}$ such that there is a book embedding of $G$ with arboricity (star-arboricity) $k$. There is no point in defining ‘book outerthickness’ since it would always equal book thickness. By definition,

$$
\begin{align*}
\theta(G) & \leq \overline{\theta}(G) \leq bt(G) \\
\theta_o(G) & \leq \overline{\theta}_o(G) \leq bt(G) \\
a(G) & \leq \overline{a}(G) \leq ba(G) \\
\overline{sa}(G) & \leq \overline{sa}(G) \leq bsa(G) .
\end{align*}
$$

5 Main Results

As summarised in Table \ref{Table1}, we determine the value of each geometric graph parameter defined in Section \ref{Section4} for $T_k$.

The following theorem is the most significant result in the paper.

**Theorem 1.** The maximum thickness and maximum geometric thickness of a graph in $T_k$ satisfy

$$
\theta(T_k) = \overline{\theta}(T_k) = [k/2] .
$$
Table 1: Maximum parameter values for graphs in $T_k$.

| type of drawing | thickness | outerthickness | arboricity | star-arboricity |
|-----------------|-----------|----------------|------------|-----------------|
| topological     | $\lceil k/2 \rceil$ | $k$           | $k$        | $k + 1$         |
| geometric       | $\lceil k/2 \rceil$ | $k$           | $k$        | $k + 1$         |
| book ($k \leq 2$) | $k$       | -              | $k + 1$    | $k + 1$         |
| book ($k \geq 3$) | $k + 1$   | -              | $k + 1$    | $k + 1$         |

Theorem 1 says that the lower bound for the thickness of $T_k$ (Proposition 1) can be matched by an upper bound, even in the more restrictive setting of geometric thickness. Theorem 1 is proved in Section 8.

**Theorem 2.** The maximum arboricity, maximum outerthickness, maximum geometric arboricity, and maximum geometric outerthickness of a graph in $T_k$ satisfy

$$a(T_k) = \theta_o(T_k) = \overline{\theta_o}(T_k) = \overline{a}(T_k) = k.$$

Theorem 2 says that our lower bounds for the arboricity and outerthickness of $T_k$ (Propositions 2 and 3) can be matched by upper bounds on the corresponding geometric parameter. By the lower bound in Proposition 3 to prove Theorem 2 it suffices to show that $\overline{a}(T_k) \leq k$; we do so in Section 8.

Now we describe our results for book embeddings.

**Theorem 3.** The maximum book thickness and maximum book arboricity of a graph in $T_k$ satisfy

$$bt(T_k) = ba(T_k) = \begin{cases} k & \text{for } k \leq 2, \\ k + 1 & \text{for } k \geq 3 \end{cases}.$$

Theorem 3 with $k = 1$ states that every tree has a 1-page book embedding, as proved by Bernhart and Kainen [13]. Rengarajan and Veni Madhavan [67] proved that every series-parallel graph has a 2-page book embedding (also see [24]); that is, $bt(T_2) \leq 2$. Note that $bt(T_2) = 2$ since there are series-parallel graphs that are not outerplanar, $K_{2,3}$ being the primary example. We prove the stronger result that $ba(T_2) = 2$ in Section 7.

Ganley and Heath [37] proved that every $k$-tree has a book embedding with thickness at most $k + 1$. In their proof, each noncrossing subgraph is in fact a star-forest. Thus

$$bt(T_k) \leq ba(T_k) \leq bsa(T_k) \leq k + 1.$$

We give an alternative proof of this result in Section 7. Ganley and Heath [37] proved a lower bound of $bt(T_k) \geq k$, and conjectured that $bt(T_k) = k$. Thus Theorem 3 refutes this conjecture. The proof is given in Section 8 where we construct a $k$-tree $Q_k$ with $bt(Q_k) \geq k + 1$. Thus Theorem 3 gives an example of an abstract parameter that is not matched by its geometric counterpart. In particular, $bt(T_k) > \theta_o(T_k) = k$ for $k \geq 3$. 
Note that Togasaki and Yamazaki [72] proved that $bt(G) \leq k$ under the stronger assumption that $G$ has pathwidth $k$. Finally observe that the lower bound in Proposition 4 and Equation 3 imply the following result.

**Corollary 1.** The maximum star-arboricity, maximum geometric star-arboricity, and maximum book star-arboricity of a graph in $T_k$ satisfy

$$sa(T_k) = bsa(T_k) = sa(T_k) = k + 1.$$ 

### 6 General Approach

When proving upper bounds, we need only consider $k$-trees, since edges can be added to a graph with treewidth $k$ to obtain a $k$-tree, without decreasing the relevant thickness or arboricity parameter. The definition of a $k$-tree $G$ suggests a natural approach to drawing $G$: choose a simplicial vertex $w$, recursively draw $G \setminus w$, and then add $w$ to the drawing. For the problems under consideration this approach fails because the neighbours of $w$ may have high degree. The following lemma solves this impasse.

**Lemma 1.** Every $k$-tree $G$ has a nonempty independent set $S$ of $k$-simplicial vertices such that either:

(a) $G \setminus S \simeq K_k$ (that is, $G \simeq K_{k,|S|}$), or

(b) $G \setminus S$ is a $k$-tree containing a $k$-simplicial vertex $v$ such that:

- for each vertex $w \in S$, there is exactly one vertex $u \in N_{G \setminus S}(v)$ such that $N_G(w) = N_{G \setminus S}[v] \setminus \{u\}$, and
- each $k$-simplicial vertex of $G$ that is not in $S$ is not adjacent to $v$.

**Proof.** Every $k$-tree has at least $k + 1$ vertices. If $|V(G)| = k + 1$ then $G \simeq K_{k+1}$ and property (a) is satisfied with $S = \{v\}$ for each vertex $v$. Now assume that $|V(G)| \geq k + 2$. Let $L$ be the set of $k$-simplicial vertices of $G$. Then $L$ is a nonempty independent set, and $G \setminus L$ is a $k$-tree or $G \setminus L \simeq K_k$. If $G \setminus L \simeq K_k$, then property (a) is satisfied with $S = L$. Otherwise, $G \setminus L$ has a $k$-simplicial vertex $v$. Let $S$ be the set of neighbours of $v$ in $L$. We claim that property (b) is satisfied. Now $S \neq \emptyset$, as otherwise $v \in L$. Since $G$ is not a clique and each vertex in $S$ is simplicial, $G \setminus S$ is a $k$-tree. Consider a vertex $w \in S$. Now $N_G(w)$ is a $k$-clique and $v \in N_G(w)$. Thus $N_G(w) \subseteq N_{G \setminus S}[v]$. Since $|N_G(w)| = k$ and $|N_{G \setminus S}[v]| = k + 1$, there is exactly one vertex $u \in N_{G \setminus S}(v)$ for which $N_G(w) = N_{G \setminus S}[v] \setminus \{u\}$. The final claim is immediate.

Lemma 1 is used to prove all of the upper bounds that follow. Our general approach is:

- in a recursively computed drawing of $G \setminus S$, draw the vertices in $S$ close to $v$,
- for each vertex $w \in S$, colour the edge $wx (x \neq v)$ by the colour assigned to $vx$, and colour the edge $wv$ by the colour assigned to the edge $vu$, where $u$ is the neighbour of $v$ that is not adjacent to $w$. 

10
Constructions of Book Embeddings

First we prove that $bsa(T_k) = k + 1$. The lower bound follows from the stronger lower bound $sa(T_k) \geq k + 1$ in Proposition 4. The upper bound is proved by induction on $|V(G)|$ with the following hypothesis. Recall that in an edge-coloured graph, a vertex $v$ is colourful if all the edges incident to $v$ receive distinct colours.

Lemma 2. Every $k$-tree $G$ has a book embedding with star-arboricity $k + 1$ such that:
- if $G \simeq K_{k+1}$ then at least one vertex is colourful, and
- if $G \not\simeq K_{k+1}$ then every $k$-simplicial vertex is colourful.

Proof. Apply Lemma 1 to $G$. We obtain a nonempty independent set $S$ of $k$-simplicial vertices of $G$.

First suppose that $G \setminus S \simeq K_k$ with $V(G \setminus S) = \{u_1, u_2, \ldots, u_k\}$. Position $V(G)$ arbitrarily on a circle, and draw the edges straight. Every edge of $G$ is incident to some $u_i$. Colour the edges $1, 2, \ldots, k$ so that every edge coloured $i$ is incident to $u_i$. Thus each colour class is a noncrossing star, and every vertex in $S$ is colourful. If $G \simeq K_{k+1}$ then $|S| = 1$ and at least one vertex is colourful. If $G \not\simeq K_{k+1}$ then no vertex $u_i$ is $k$-simplicial; thus every $k$-simplicial vertex is in $S$ and is colourful.

Otherwise, by Lemma 1(b), $G \setminus S$ is a $k$-tree containing a $k$-simplicial vertex $v$, such that $N_{G\setminus S}(v)$ for each vertex $w \in S$. Say $N_{G\setminus S}(v) = \{u_1, u_2, \ldots, u_k\}$.

Apply the induction hypothesis to $G \setminus S$. If $G \setminus S \simeq K_{k+1}$ then we can nominate $v$ to be a vertex of $G \setminus S$ that becomes colourful. By induction, we obtain a book embedding of $G \setminus S$ with star-arboricity $k + 1$, in which $v$ is colourful. Without loss of generality, each edge $vu_i$ is coloured $i$. Let $x$ be a vertex next to $v$ on the convex hull. Position the vertices in $S$ arbitrarily between $v$ and $x$. For each $w \in S$, colour each edge $wu_i$ by $i$, and colour $wv$ by $k + 1$, as illustrated in Figure 2(a).

By construction, each vertex in $S$ is colourful. The edges $\{vw : w \in S\}$ form a new star component of the star-forest coloured $k + 1$. For each colour $i \in \{1, 2, \ldots, k\}$, the component
of the subgraph of $G \setminus S$ that is coloured $i$ and contains $v$ is a star rooted at $u_i$ with $v$ a leaf. Thus it remains a star by adding the edge $wu_i$ for all $w \in S$.

Suppose that two edges $e$ and $f$ of $G$ cross and are both coloured $i$ ($i \in \{1, 2, \ldots, k\}$). Then $e$ and $f$ are not both in $G \setminus S$. Without loss of generality, $e$ is incident to a vertex $w \in S$. The edges of $G$ that are coloured $i$ and have at least one endpoint in $S \cup \{v\}$ form a noncrossing star (rooted at $u_i$ if $1 \leq i \leq k$, and rooted at $v$ if $i = k + 1$). Thus $f$ has no endpoint in $S \cup \{v\}$. Observe that $vw$ crosses no edge in $G \setminus S$. Thus $e = wu_i$. Since $S \cup \{v\}$ is consecutive on the circle and $f$ has no endpoint in $S \cup \{v\}$, $f$ also crosses $vu_i$. Hence $f$ and $vu_i$ are two edges of $G \setminus S$ that cross and are both coloured $i$. This contradiction proves that no two edges of $G$ cross and receive the same colour.

It remains to prove that every $k$-simplicial vertex in $G$ is colourful. Each vertex in $S$ is colourful. Consider a $k$-simplicial vertex $x$ of $G$ that is not in $S$. By Lemma 1(b), $x$ is not adjacent to $v$. Thus $x$ is adjacent to no vertex in $S$, and $x$ is $k$-simplicial in $G \setminus S$. Moreover, $G \setminus S$ is not complete. By induction, $x$ is colourful in $G \setminus S$ and in $G$.

Now we prove Theorem 3 with $k = 2$, which states that $bt(T_2) = ba(T_2) = 2$. The lower bound $bt(T_2) \geq 2$ holds since $K_{2,3}$ is series-parallel but is not outerplanar. We prove the upper bound $ba(T_2) \leq 2$ by induction on $|V(G)|$ with the following hypothesis.

**Lemma 3.** Every 2-tree $G$ has a book embedding with arboricity 2 such that:

- if $G \simeq K_3$ then two vertices are colourful, and
- if $G \not\simeq K_3$ then every 2-simplicial vertex is colourful.

**Proof.** Apply Lemma 1 to $G$. We obtain a nonempty independent set $S$ of 2-simplicial vertices of $G$.

First suppose that $G \setminus S \simeq K_2$ with $V(G \setminus S) = \{u_1, u_2\}$. Position $V(G)$ at distinct points on a circle in the plane, and draw the edges straight. Every edge is incident to $u_1$ or $u_2$. Colour every edge incident to $u_1$ by 1. Colour every edge incident to $u_2$ (except $u_1u_2$) by 2. Thus each colour class is a noncrossing star, and each vertex in $S$ is colourful. If $G \simeq K_3$ then $|S| = 1$ and $u_2$ is also colourful. If $G \not\simeq K_3$ then neither $u_1$ nor $u_2$ are 2-simplicial; thus each 2-simplicial vertex is colourful.

Otherwise, by Lemma 1(b), $G \setminus S$ is a 2-tree containing a 2-simplicial vertex $v$. Say $N_{G \setminus S}(v) = \{u_1, u_2\}$. For every vertex $w \in S$, $N_G(w) = \{v, u_1\}$ or $N_G(w) = \{v, u_2\}$. Let $S_1 = \{w \in S : N_G(w) = \{v, u_1\}\}$ and $S_2 = \{w \in S : N_G(w) = \{v, u_2\}\}$.

Apply the induction hypothesis to $G \setminus S$. If $G \setminus S \simeq K_3$ we can nominate $v$ to be a vertex of $G \setminus S$ that becomes colourful. By induction, we obtain a book embedding of $G \setminus S$ with arboricity 2, in which $v$ is colourful. Without loss of generality, each edge $vu_i$ is coloured $i$. Say $u_1$ appears before $u_2$ in clockwise order from $v$. Say $(x, v, y)$ are consecutive in clockwise order, as illustrated in Figure 2(b). Position the vertices in $S_1$ between $v$ and $y$, and position the vertices in $S_2$ between $x$ and $v$. For all $w \in S_i$, colour each edge $wu_i$ by $i$, and colour $vw$ by $3 - i$.

The only edge that can cross an edge $wv$ ($w \in S_i$) is some $pu_i$ where $p \in S_i$. These edges receive distinct colours. If an edge $e$ of $G \setminus S$ crosses some edge $wu_i$, then $e$ also crosses
vu_i (since deg_{G \setminus S}(v) = 2). Since wu_i receives the same colour as vu_i, e must be coloured differently from wu_i. Hence edges assigned the same colour do not cross.

By construction, each vertex w ∈ S is colourful; w becomes a leaf in both forests of the partition. It remains to prove that every 2-simplicial vertex in G is colourful. Each vertex in S is colourful. Consider a k-simplicial vertex x of G that is not in S. By Lemma 1(b), x is not adjacent to v. Thus x is adjacent to no vertex in S, and x is 2-simplicial in G \ S. Moreover, G \ S is not complete. By induction, x is colourful in G \ S and in G. □

8 Constructions of Geometric Drawings

In this section we prove Theorems 1 and 2. First we introduce some geometric notation. Let v and w be distinct points in the plane; see Figure 3. Let \( \overrightarrow{vw} \) be the line through v and w. Let \( \overline{vw} \) be the open line-segment with endpoints v and w. Let \( \overline{\overline{vw}} \) be the closed line-segment with endpoints v and w. Let \( \overrightarrow{\overline{vw}} \) be the open ray from v through w. Let \( \overleftarrow{vw} \) be the open ray opposite to \( \overrightarrow{\overline{vw}} \); that is, \( \overleftarrow{vw} := (\overrightarrow{\overline{vw}} \setminus \overrightarrow{vw}) \setminus \{v\} \).

![Figure 3: Notation for lines and rays.](image)

For every point p ∈ \( \mathbb{R}^2 \) and set of points Q ⊂ \( \mathbb{R}^2 \setminus \{p\} \), such that \( Q \cup \{p\} \) is in general position, let

\[ R(p, Q) := \{pq, qp : q ∈ Q\} \]

be the set of rays from p to the points in Q together with their opposite rays, in clockwise order around p. (Since \( Q \cup \{p\} \) is in general position, the rays in \( R(p, Q) \) are pairwise disjoint, and their clockwise order is unique.)

Let r and r’ be non-collinear rays from a single point v. The wedge \( \triangledown(r, r') \) centred at v is the unbounded region of the plane obtained by sweeping a ray from r to r’ through the lesser of the two angles formed by r and r’ at v. We consider \( \triangledown(r, r') \) to be open in the sense that \( r \cup r' \cup \{v\} \) does not intersect \( \triangledown(r, r') \).

The proofs of Theorems 1 and 2 are incremental constructions of geometric drawings. The insertion of new vertices is based on the following definitions.

Consider a geometric drawing of a graph G. Let v be a vertex of G. For \( \varepsilon > 0 \), let \( D_\varepsilon(v) \) be the open disc of radius \( \varepsilon \) centred at v. For a point u, let

\[ C_\varepsilon(v, u) := \bigcup \{\overline{uv} : x ∈ D_\varepsilon(v)\} \]

be the region in the plane consisting of the union of all open line-segments from u to the points in \( D_\varepsilon(v) \). Let

\[ T_\varepsilon(v) := \bigcup \{C_\varepsilon(v, u) : u ∈ N_G(v)\} \]

be the region in the plane consisting of the union of all open line-segments from each neighbour of v to the points in \( D_\varepsilon(v) \).
As illustrated in Figure 4(a), a vertex \( v \) in a general position geometric drawing of a graph \( G \) is \( \varepsilon \)-empty if:
(a) the only vertex of \( G \) in \( T_\varepsilon(v) \) is \( v \),
(b) every edge of \( G \) that intersects \( D_\varepsilon(v) \) is incident to \( v \),
(c) \( (V(G) \setminus \{v\}) \cup \{p\} \) is in general position for each point \( p \in D_\varepsilon(v) \), and
(d) the clockwise orders of \( R(v, N_G(v)) \) and \( R(p, N_G(v)) \) are the same for each point \( p \in D_\varepsilon(v) \).

![Figure 4:](image)

**Observation 1.** Every vertex \( v \) in a general position geometric drawing of a graph \( G \) is \( \varepsilon \)-empty for some \( \varepsilon > 0 \).

**Proof.** Consider the arrangement \( A \) consisting of the lines through every pair of vertices in \( G \setminus v \); see [57] for background on line arrangements. Since \( V(G) \) is in general position, \( v \) is in some cell \( C \) of \( A \). Since \( C \) is an open set, there exists \( \varepsilon > 0 \) such that \( D_\varepsilon(v) \subset C \). For every neighbour \( u \in N_G(v) \), no vertex \( x \) of \( G \setminus v \) is in \( C_\varepsilon(v, u) \), as otherwise \( xu \) would intersect \( D_\varepsilon(v) \). Thus property (a) holds. No line of \( A \) intersects \( C \). In particular, no edge of \( G \setminus v \) intersects \( C \), and property (b) holds. No point \( p \in D_\varepsilon(v) \) is collinear with two vertices of \( G \setminus v \), as otherwise \( D_\varepsilon(v) \) would intersect a line in \( A \). Thus property (c) holds. The radial order of \( V(G) \setminus v \) is the same from each point in \( C \). In particular, property (d) holds. Therefore \( v \) is \( \varepsilon \)-empty. \( \square \)
Observation 2. Let \( v \) be an \( \varepsilon \)-empty vertex in a general position geometric drawing of a graph \( G \). Let \( u \in N_G(v) \). Suppose that some edge \( e \in E(G) \) crosses \( \overline{vu} \) for some point \( p \in D_\varepsilon(v) \). Then either \( e \) is incident to \( v \), or \( e \) also crosses the edge \( vu \).

Proof. If \( e \) is incident to \( v \), then we are done. Now assume that \( e \) is not incident to \( v \). Thus \( e \) does not intersect \( D_\varepsilon(v) \) by property (b) of the choice of \( \varepsilon \). Since \( p \in D_\varepsilon(v) \), we have \( \overline{vu} \subset C_\varepsilon(v,u) \). Thus the crossing point between \( e \) and \( \overline{vu} \) is in \( C_\varepsilon(v,u) \setminus D_\varepsilon(v) \). In particular, \( e \) intersects \( C_\varepsilon(v,u) \). By property (a) of the choice of \( \varepsilon \) and since \( e \) is not incident to \( v \), no endpoint of \( e \) is in \( T_\varepsilon(v) \).

We have proved that \( e \) does not intersect \( D_\varepsilon(v) \), \( e \) intersects \( C_\varepsilon(v,u) \), and no endpoint of \( e \) is in \( T_i(v) \). Observe that any segment with these three properties must cross \( vu \). Thus \( e \) crosses \( vu \).

\[ \square \]

8.1 Proof of Theorem 2

Theorem 2 states that \( a(T_k) = \theta_0(T_k) = \theta_\alpha(T_k) = \overline{\alpha}(T_k) = k \). By the discussion in Section 5, it suffices to show that for geometric arboricity, \( \overline{\alpha}(T_k) \leq k \). We proceed by induction on \( |V(G)| \) with the following hypothesis.

Proposition 5. Every \( k \)-tree \( G \) has a general position geometric drawing with arboricity \( k \) such that:

- if \( G \cong K_{k+1} \) then at least one vertex is colourful, and
- if \( G \not\cong K_{k+1} \) then every \( k \)-simplicial vertex is colourful.

Proof. Apply Lemma 1 to \( G \). We obtain a nonempty independent set \( S \) of \( k \)-simplicial vertices of \( G \).

First suppose that \( G \setminus S \cong K_k \) with \( V(G \setminus S) = \{u_1,u_2,\ldots,u_k\} \). Fix an arbitrary general position geometric drawing of \( G \). Greedily colour the edges of \( G \) with colours \( 1,2,\ldots,k \), starting with the edges incident to \( u_1 \) and ending with the edges incident to \( u_k \), so that every edge coloured \( i \) is incident to \( u_i \). Thus each colour class is a noncrossing star, and every vertex in \( S \) is colourful. If \( G \cong K_{k+1} \) then \(|S| = 1 \) and at least one vertex is colourful. If \( G \not\cong K_{k+1} \) then no vertex \( u_i \) is \( k \)-simplicial in \( G \); thus each \( k \)-simplicial vertex is in \( S \) and is colourful.

Otherwise, by Lemma 1(b), \( G \setminus S \) is a \( k \)-tree containing a \( k \)-simplicial vertex \( v \). Say \( N_{G \setminus S}(v) = \{u_1,u_2,\ldots,u_k\} \). Each vertex \( w \in S \) has \( N_G(w) = N_{G \setminus S}[v] \setminus \{u_i\} \) for exactly one value of \( i \in \{1,2,\ldots,k\} \). Let \( S_i := \{w \in S : N_G(w) = N_{G \setminus S}[v] \setminus \{u_i\} \} \) for each \( i \in \{1,2,\ldots,k\} \). Then \( \{S_1,S_2,\ldots,S_k\} \) is a partition of \( S \).

Apply the induction hypothesis to \( G \setminus S \). If \( G \setminus S \cong K_{k+1} \) then we can nominate \( v \) to be a vertex of \( G \setminus S \) that becomes colourful. By induction, we obtain a general position geometric drawing of \( G \setminus S \) with arboricity \( k \), in which \( v \) is colourful. Without loss of generality, each edge \( vu_i \) is coloured \( i \).

By Observation 1, \( v \) is \( \varepsilon \)-empty in the general position geometric drawing of \( G \setminus S \) for some \( \varepsilon > 0 \). Let \( X_1,X_2,\ldots,X_k \) be pairwise disjoint wedges centred at \( v \) such that \( \overline{vu_i} \subset X_i \) for all \( i \in \{1,2,\ldots,k\} \). Position the vertices of \( S_i \) in \( X_i \cap D_\varepsilon(v) \) so that \( V(G) \) is in general position. This is possible since \( X_i \cap D_\varepsilon(v) \) is an open (infinite) region, but there are only finitely many pairs of vertices. Draw each edge straight. For each vertex \( w \in S_i \), colour the edge \( uv \) by \( i \).
and colour the edge \(wu_j\) \((j \neq i)\) by \(j\). Thus \(w\) is colourful; \(w\) becomes a leaf in each of the \(k\) forests. This construction is illustrated in Figure 4(b).

To prove that edges assigned the same colour do not cross, consider the set of edges coloured \(i\) to be partitioned into three sets:
\(\begin{align*}
(1) \text{edges in } G \setminus S \text{ that are coloured } i, \\
(2) \text{edges } wu_i \text{ for some } w \in S \setminus S_i, \text{ and} \\
(3) \text{edges } vw \text{ for some } w \in S_i.
\end{align*}\)

Type-(1) edges do not cross by induction. Type-(2) edges do not cross since they are all incident to \(u_i\). Type-(3) edges do not cross since they are all incident to \(v\).

Suppose that a type-(1) edge \(e\) crosses a type-(2) edge \(wu_i\) for some \(w \in S\). By Observation 2 with \(p = w \in D_{\epsilon}(v)\), either \(e\) is incident to \(v\), or \(e\) also crosses \(vu_i\). Since \(e\) and \(vu_i\) are both coloured \(i\), they do not cross in \(G\), and we can now assume that \(e\) is incident to \(v\). Thus \(e = vu_i\), which is the only edge in \(G \setminus S\) that is incident to \(v\) and is coloured \(i\). Since \(e\) and \(wu_i\) have a common endpoint, \(e\) and \(wu_i\) do not cross, which is a contradiction. Thus type-(1) and type-(2) edges do not cross.

Now suppose that a type-(1) edge \(e\) crosses a type-(3) edge \(vw\) for some \(w \in S_i\). Then \(e \neq vu_i\), since \(vu_i\) and \(vw\) have a common endpoint. Now, \(vw\) is contained in \(D_{\epsilon}(v)\). Thus \(e\) intersects \(D_{\epsilon}(v)\), which contradicts property (b) of the choice of \(\epsilon\). Thus type-(1) and type-(3)
edges do not cross.

By construction, no type-(2) edge intersects the wedge \(X_i\). Since every type-(3) edge is contained in \(X_i\), type-(2) and type-(3) edges do not cross. Therefore edges assigned the same colour do not cross.

It remains to prove that each \(k\)-simplicial vertex of \(G\) is colourful. Each vertex in \(S\) is colourful. Consider a \(k\)-simplicial vertex \(x\) that is not in \(S\). By Lemma 1b, \(x\) is not adjacent to \(v\). Thus \(x\) is adjacent to no vertex in \(S\), and \(x\) is \(k\)-simplicial in \(G \setminus S\). Moreover, \(G \setminus S\) is not complete. By induction, \(x\) is colourful in \(G \setminus S\) and in \(G\).

\[\square\]

8.2 Proof of Theorem 1

Theorem 1 states that \(\theta(T_k) = \overline{\theta}(T_k) = \lfloor k/2 \rfloor\). The thickness lower bound, \(\theta(T_k) \geq \lfloor k/2 \rfloor\), is Proposition 1. For the upper bound on the geometric thickness, \(\overline{\theta}(T_k) \leq \lfloor k/2 \rfloor\), it suffices to prove that \(\overline{\theta}(T_{2k}) \leq k\) for all \(k \geq 2\) (since graphs in \(T_2\) are planar, and thus have geometric thickness 1). We use the following definitions, for some fixed \(k \geq 2\).

Let
\[I := \{i, -i : 1 \leq i \leq k\}\]

Suppose that \(\phi\) is a geometric drawing of a graph \(G\). (Note that \(G\) is not necessarily a \(2k\)-tree, and \(\phi\) is not necessarily in general position.) Suppose that \(v\) is a vertex of \(G\) with degree \(2k\), where
\[N_G(v) = (u_1, u_2, \ldots, u_k, u_{k-1}, u_{k-2}, \ldots, u_1)\] (4)
are the neighbours of \(v\) in clockwise order around \(v\) in \(\phi\). (Since no edge passes through a vertex, this cyclic ordering is well defined.) For each \(i \in I\), define the \(i\)-\(wedge\) of \(v\) (with respect to the labelling of \(N_G(v)\) in Equation (4)) to be
\[F_i(v) := \nabla(\overleftarrow{vu_i}, \overrightarrow{vu_{-i}})\].

16
If \( u_i, v, u_j \) are collinear, then \( \overrightarrow{vu_i} = \overrightarrow{vu_j} \). But if \( \phi \) is in general position, then \( \overrightarrow{vu_i} \neq \overrightarrow{vu_j} \) for all \( i, j \in I \). Now suppose that, in addition, \( \phi \) is in general position. Let

\[
R(v) := R(v, N_G(v)) = \{ \overrightarrow{vu_i}, \overrightarrow{vu_j} : i \in I \}
\]

be the set of \( 2k \) open rays from \( v \) through its neighbours together with their \( 2k \) opposite open rays, in clockwise order around \( v \) in \( \phi \). We say \( v \) is balanced in \( \phi \) if \( \overrightarrow{vu_i} \) and \( \overrightarrow{vu_{i-1}} \) are consecutive in \( R(v) \) for each \( i \in I \). Note that \( v \) is balanced if and only if \( F_i(v) \cap F_j(v) = \emptyset \) for all distinct \( i, j \in I \). Moreover, whether \( v \) is balanced does not depend on the choice of labelling in Equation (4).

Now suppose that, in addition, \( G \) is a \( 2k \)-tree, and \( \phi \) has thickness \( k \). Consider the edges of \( G \) to be coloured 1, 2, \ldots, \( k \), where edges of the same colour do not cross in \( \phi \). As illustrated in Figure 5(a), a \( 2k \)-simplicial vertex \( v \) of \( G \) is a fan in \( \phi \) if, for some labelling of \( N_G(v) \) as in Equation (4), we have:

- \( v \) is balanced in \( \phi \), and
- the edge \( vu_i \) is coloured \( |i| \) for each \( i \in I \).

Note that for all \( Q \subseteq V(G) \) and all \( v \in Q \) such that \( G[Q] \) is a \( 2k \)-tree and \( v \) is \( 2k \)-simplicial in \( G \), \( v \) is a fan in \( \phi \) if and only if \( v \) is a fan in the drawing of \( G[Q] \) induced by \( \phi \).

A drawing \( \phi \) of a \( 2k \)-tree \( G \) is good if:

- \( \phi \) is a general position geometric drawing,
- \( \phi \) has thickness \( k \),
- if \( G \simeq K_{2k+1} \) then at least one vertex of \( G \) is a fan in \( \phi \), and
- if \( G \not\simeq K_{2k+1} \) then every \( 2k \)-simplicial vertex of \( G \) is a fan in \( \phi \).

The proof of Theorem 1 uses the following two lemmas about constructing good drawings.

---

4Since \( G[Q] \) is a \( 2k \)-tree, it has minimum degree \( 2k \). Since \( v \in Q \) and \( \deg_G(v) = 2k \), we have \( \deg_{G[Q]}(v) = 2k \). Thus every neighbour of \( v \) in \( G \) is also in \( Q \). Thus \( v \) is \( 2k \)-simplicial in \( G[Q] \).
Lemma 4. Consider a 2k-tree $G$ for some $k \geq 2$. Suppose that $G$ has a good drawing $\phi$, and $v$ is a fan vertex in $\phi$. Let $G'$ be the 2k-tree obtained from $G$ by adding a new vertex $w$ onto $N_G(v)$. Then $w$ can be inserted into $\phi$ to obtain a good drawing $\phi'$ of $G'$.

Proof. Say $(u_1, u_2, \ldots, u_k, u_{-1}, u_{-2}, \ldots, u_{-k})$ are the neighbours of $v$ in clockwise order around $v$. Since $v$ is a fan in $\phi$, the edge $vu_i$ is coloured $|i|$ for all $i \in I$. By Observation 1, $v$ is $\varepsilon$-empty for some $\varepsilon > 0$. Let

$$X := D_{\varepsilon}(v) \setminus \{v\} \cup \{F_i(v) : i \in I\}.$$ 

Thus $X$ consists of 2k connected sets having nonempty interior. Hence, there is a nonempty, in fact open, subset of $X$ consisting of points that are not collinear with any two distinct vertices of $G$. Map $w$ to any point in that subset, and draw each edge $wu_i$ straight $(i \in I)$. We obtain a general position geometric drawing $\phi'$ of $G'$. As illustrated in Figure 5(b), colour each edge $wu_i$ of $G'$ by $|i|$, which is the same colour assigned to $vu_i$.

Consider an edge $e$ of $G$ that crosses $wu_i$ in $\phi'$ for some $i \in I$. By construction, $wu_i$ is coloured $|i|$. Suppose, for the sake of contradiction, that $e$ is also coloured $|i|$. By Observation 2 with $p = w \in (D_{\varepsilon}(v))$, either $e$ is incident to $v$, or $e$ also crosses $vu_i$. Since $e$ and $vu_i$ are both coloured $i$ in $G$, $e$ does not cross $vu_i$, and we can now assume that $e$ is incident to $v$. Since $vu_i$ and $wu_i$ share an endpoint, $e \neq vu_i$. Thus $e = vu_{-i}$, which is the only other edge incident to $v$ coloured $|i|$. Since $wu_i$ crosses $vu_{-i}$, we have that $w \in F_{-i}(v)$, which contradicts the placement of $w$. Thus edges of $G'$ that are assigned the same colour do not cross in $\phi'$.

Let $x \neq w$ be a 2k-simplicial vertex in $G'$. Then $x$ is not adjacent to $w$, and $x$ is 2k-simplicial in $G$. Since $x$ is a fan in $\phi$, it also is a fan in $\phi'$. We now prove that $w$ is a fan in $\phi'$. By property (d) of the choice of $\varepsilon$, and since $w \in D_{\varepsilon}(v)$, the cyclic orderings of the ray sets $R(v)$ and $R(w)$ are the same. Since $v$ is a fan in $\phi$, and by the colouring of the edges incident to $w$, $w$ is also fan in $\phi'$.

If $G \simeq K_{2k+1}$, then $v$ and $w$ are the only 2k-simplicial vertices in $G'$, and thus every 2k-simplicial vertex of $G'$ is a fan in $\phi'$. If $G \not\simeq K_{2k+1}$, consider a 2k-simplicial vertex $y \neq w$ of $G'$. No pair of 2k simplicial vertices in $G'$ are adjacent. Thus $y$ is 2k-simplicial in $G$ and $y$ is a fan in $\phi$ (and $\phi'$). Thus every 2k-simplicial vertex of $G'$ is a fan in $\phi'$, as required.

Lemma 5. For all $k \geq 2$, the complete graph $K_{2k+1}$ has a good drawing in which any given vertex $v$ is a fan.

Proof. Say $V(K_{2k+1}) = \{v, u_1, u_2, \ldots, u_{2k}\}$. As illustrated in Figure 6(a), position $u_1, u_2, \ldots, u_{2k}$ evenly spaced and in this order on a circle in the plane centred at a point $p$. The edges induced by $\{u_1, u_2, \ldots, u_{2k}\}$ can be $k$-coloured using the standard book embedding of $K_{2k}$ with thickness $k$: colour each edge $u_\alpha u_\beta$ by $1 + \lfloor (\alpha + \beta) \mod 2k \rfloor$. Then the colours are $1, 2, \ldots, k$, and each colour class forms a noncrossing zig-zag subgraph $\mathcal{Z}$. [13, 28]

Rename each vertex $u_{k+i}$ by $u_{-i}$. As illustrated in Figure 6(b), the edges $\{u_i u_{-i} : 1 \leq i \leq k\}$ pairwise intersect at $p$. Position $v$ strictly inside a cell of the drawing of $K_{2k}$ that borders $p$ (the shaded region in Figure 6(a)). Then $V(K_{2k+1})$ is in general position. For all $i \in I$, colour $vu_i$ by $|i|$. Then edges assigned the same colour do not cross. $v$ is a fan since $R(v, \{u_i : i \in I\}) = (\overline{vu_{-1}}, \overline{vu_1}, \overline{vu_{-2}}, \overline{vu_2}, \ldots, \overline{vu_{-k}}, \overline{vu_k})$. \qed
Proposition 6. For all $k \geq 2$, every $2k$-tree $G$ has a good drawing.

Proof. In this proof we repeatedly use two indices, $i$ and $r$, whose ranges remain unchanged; in particular, $i \in I$ and $r \in \{1, 2, \ldots, k\}$.

We proceed by induction on $|V(G)|$. If $G \cong K_{2k+1}$ the result is Lemma 5. Now assume that $G \not\cong K_{2k+1}$. Apply Lemma 1 to $G$. We obtain a nonempty independent set $S$ of $2k$-simplicial vertices of $G$.

First suppose that $G \setminus S \cong K_{2k}$. Let $v$ be an arbitrary vertex in $S$. By Lemma 5, $G \setminus (S \setminus \{v\}) (\cong K_{2k+1})$ has a good drawing in which $v$ is a fan. By Lemma 4, each vertex $w \in S \setminus \{v\}$ can be inserted into the drawing (one at the time) resulting in a good drawing of $G$.

Otherwise, by Lemma 1(b), $G \setminus S$ is a $2k$-tree containing a $2k$-simplicial vertex $v$, such that $N_{G \setminus S}(v) \subset N_{G \setminus S}[v]$ for each vertex $w \in S$.

Apply the induction hypothesis to $G \setminus S$. If $G \setminus S \cong K_{2k+1}$ then we can nominate $v$ to be a vertex of $G \setminus S$ that is a fan. By induction, we obtain a good drawing $\phi$ of $G \setminus S$ in which $v$ is a fan. Say $N_{G \setminus S}(v) = (u_1, u_2, \ldots, u_k, u_{-1}, u_{-2}, \ldots, u_{-k})$ in clockwise order about $v$. Thus each edge $vu_i$ is coloured $|i|$.

By Lemma 1(b), for each vertex $w \in S$, there is exactly one $i \in I$ for which $N_G(w) = N_{G \setminus S}[v] \setminus \{u_i\}$. Let $S_i := \{w \in S : N_G(w) = N_{G \setminus S}[v] \setminus \{u_i\}\}$ for each $i \in I$. The vertices in $S_i$ have the same neighbourhood in $G$, and $\{S_i : i \in I\}$ is a partition of $S$.

For each $i \in I$, choose one vertex $x_i \in S_i$ (if any). Let $Q := \{x_i : i \in I\}$. Suppose we have a good drawing of $(G \setminus S) \cup Q$. Then by Lemma 1, each vertex $w \in S \setminus Q$ can be inserted into the drawing (one at the time) resulting in a good drawing of $G$. Thus, from now on, we can assume that $S = Q (= \{x_i : i \in I\})$. Below we describe how to insert the vertices $\{x_i : i \in I\}$ into $\phi$ to obtain a good drawing $\phi'$ of $G$.

First we colour the edges incident to each vertex $x_i \in S$. Colour $x_iv$ by $|i|$, and colour

Figure 6: (a) Book embedding of $K_{2k}$. (b) Geometric drawing of $K_{2k+1}$ in which $v$ is a fan.
\[ x_i u_j \text{ by } |j| \text{ for all } j \in I \setminus \{i\}. \] Thus there are exactly two edges of each colour incident to \( x_i \). In particular, \( x_i v \) and \( x_i u_{-i} \) are coloured \(|i|\), and \( x_i u_j \) and \( x_i u_{-j} \) are coloured \( j \) for each \( j \in \{1, 2, \ldots, k\} \setminus \{|i|\}. \)

For each \( r \in \{1, 2, \ldots, k\} \), let \( G_r \) be the spanning subgraph of \( G \) consisting of all the edges of \( G \) coloured \( r \). Let \( G^* \) be the spanning subgraph of \( G \) with edge set \( E(G) \setminus \{vu_i : i \in I\} \). Let \( G^*_r := G_r \cap G^* \) for each \( r \in \{1, 2, \ldots, k\} \).

As an intermediate step, we now construct a geometric drawing \( \phi^* \) of \( G^* \) (not in general position), in which each subgraph \( G^*_r \) is noncrossing. We later modify \( \phi^* \), by moving each vertex \( x_i \) and drawing each edge \( vu_i \), to obtain a general position geometric drawing \( \phi' \) of \( G \), in which each subgraph \( G_r \) is noncrossing.

First, let \( \phi^*(w) := \phi(w) \) for every vertex \( w \) of \( G \setminus S \). By Observation 1, \( v \) is \( \varepsilon \)-empty in \( \phi \) for some \( \varepsilon > 0 \). We now position each vertex \( x_i \) on the segment \( \overline{vu_i} \cap D_\varepsilon(v) \). We have \( F_j(x_i) = \bigvee (x_i u_j, \overline{x_i u_{-j}}) \) for all \( j \in I \setminus \{i, -i\} \). Observe that with \( x_i \in \overline{vu_i} \cap D_\varepsilon(v) \), we have \( v \not\in F_j(x_i) \) for all \( j \in I \setminus \{i, -i\} \). Therefore, for \( i \in I \) in some arbitrary order, each vertex \( x_i \) can be positioned on the segment \( \overline{vu_i} \cap D_\varepsilon(v) \) so that:

\[ \begin{itemize}
  \item \( x_i \not\in F_j(x_i) \) for each \( \ell \in I \setminus \{i\} \) and \( j \in I \setminus \{\ell, -\ell\} \), and
  \item \( V(G) \) is in general position except for the collinear triples \( v, x_i, u_{-i} \) (\( i \in I \)).
\end{itemize} \]

This is possible by the previous observation, since there is always a point close enough to \( v \) where \( x_i \) can be positioned. This placement of vertices of \( G^* \) determines a geometric drawing \( \phi^* \) of \( G^* \). The construction is illustrated in Figure 7.

**Claim 1.** The subgraph \( G^*_r \) is noncrossing in \( \phi^* \) for each \( r \in \{1, 2, \ldots, k\} \).

**Proof.** Distinguish the following three types of edges in \( G^*_r \):

1. edges of \( G^*_r \setminus S \),
2. edges \( x_r v, x_r u_{-r}, x_{-r} v, \) and \( x_{-r} u_i \),
3. edges \( x_j u_r \) and \( x_l u_{-r} \) for distinct \( j, l \in I \setminus \{r, -r\} \).

   First note that \( x_r v \cup x_r u_{-r} = vu_r \), and similarly \( x_{-r} v \cup x_{-r} u_r = vu_r \). Since no pair of edges in \( \{vu_r, vu_r\} \cup E(G_r \setminus S) \) cross in \( \phi \), no pair of edges in \( \{x_r v, x_r u_{-r}, x_{-r} v, x_{-r} u_r\} \cup E(G_r \setminus S) \) cross in \( \phi^* \).

   It remains to prove that no type-(1) edge crosses a type-(3) edge, no type-(2) edge crosses a type-(3) edge, and that no pair of type-(3) edges cross in \( \phi^* \).

20
Figure 7: Placing each $x_i$ on the segment $vu_{-i}$; intuitively speaking, the circle $D_\varepsilon$ is chosen small enough so that the edges incident with $u_i$ are almost parallel.

Consider a type-(1) edge $e$ and a type-(3) edge. Since $v$ is a fan in $\phi$, the only two edges
Claim 2. Proof. Let $S$ by induction, subgraph $G$ simplicial vertex $y$ in position. Suppose that an edge $e$ and $x$ adjacent to no vertex in $S$, $\phi$ $G$ edges of $\phi$. We obtain a number $\epsilon_r$ such that if $\phi'(v) \in D_{\epsilon_r}(w)$ for every vertex $v \in V(G^*_r)$, then $\phi'$ is a noncrossing geometric drawing of $G^*_r$ with the property that if three vertices $\phi'(a), \phi'(b), \phi'(c)$ are collinear in $\phi'$, then $\phi^*(a), \phi^*(b), \phi^*(c)$ are collinear in $\phi^*$.

Let $\delta := \min \{ \varepsilon_r : r \in \{1, 2, \ldots, k\} \}$. For each $i \in I$, let $\phi'(x_i)$ be some point in the region $D_{\delta}(x_i) \cap F_{-i}(v)$. Let $\phi'(w) := \phi(w)$ for every other vertex of $G$.

We now prove that each subgraph $G_r$ is noncrossing in $\phi'$. By Lemma $\textbf{8}$ since $\delta \leq \varepsilon_r$, each subgraph $G^*_r$ is noncrossing in $\phi'$. We must also show that the edges $vu_r$ and $vu_{-r}$ do not cross any edge in $G_r$. First note that $vu_r$ and $vu_{-r}$ do not cross since they have a common endpoint. Suppose that an edge $e$ of $G^*_r$ crosses $vu_{-r}$. Since the interior of the triangle $vx_r, u_{-r}$ contains no vertex, $e$ also crosses $vx$ or $x_r u_{-r}$. This is impossible, since $vx_r$ and $x_r u_{-r}$ are edges of $G^*_r$. Similarly, an edge $e$ of $G^*_r$ does not cross $vu_r$. Thus $G_r$ is noncrossing.

We now prove that $\phi'$ is in general position. By Lemma $\textbf{8}$ if three vertices are collinear in $\phi'$ then they are collinear in $\phi^*$. The only collinear triples in $\phi^*$ are $v, x_i, u_{-i}$ for $i \in I$. Since $\phi'(x_i)$ is in (the interior of) $F_{-i}(v)$, the vertices $v, x_i, u_{-i}$ are not collinear in $\phi'$. Thus $\phi'$ is in general position.

It remains to prove that every $2k$-simplicial vertex of $G$ is a fan in $\phi'$. Consider a $2k$-simplicial vertex $y$ that is not in $S$. By Lemma $\textbf{11}$, $y$ is not adjacent to $v$. Thus $y$ is adjacent to no vertex in $S$, and $y$ is $2k$-simplicial in $G \setminus S$. Moreover, $G \setminus S$ is not complete. By induction, $y$ is a fan in the drawing of $G \setminus S$ induced by $\phi'$, and thus $y$ is a fan in $\phi'$. Each vertex in $S$ is a fan in $\phi'$ by the following claim.

Claim 2. For each $i \in I$, the vertex $x_i$ is a fan in $\phi'$.

Proof. Let $H$ be the $2k$-tree obtained from $G \setminus S$ by adding a new vertex $h$ onto the $2k$-clique
$N_{G \setminus S}(v) = \{u_i, u_{-i} : i \in I\}$. Consider the general position geometric drawing $\sigma$ of $H$ induced by $\phi$ with $\sigma(h) := \phi'(x_i)$.

By construction, $\sigma(h) \in D_\varepsilon(v) \cap F_{-i}(v)$. Thus property (d) of the choice of $\varepsilon$ implies that the clockwise orders of $R(v, N_{G \setminus S}(v))$ and $R(h, N_H(h))$ are the same. Since $v$ is balanced in $\phi$, $h$ is balanced in $\sigma$.

Now consider the drawings $\phi'$ of $G$ and $\sigma$ of $H$. Note that $F_j(x_i) = F_j(h)$ for all $j \in I \setminus \{i, -i\}$. Furthermore, since $F_{-i} \subseteq F_i(h) \cup F_{-i}(h)$, we have $F_i(x_i) \subseteq F_i(h)$ and $F_{-i}(x_i) \subseteq F_{-i}(h)$. Therefore, $x_i$ is balanced in $\phi'$. The edges $x_i v$ and $x_i u_{-i}$ are both coloured $|i|$, and $x_i u_j$ is coloured $|j|$ for all $j \in I \setminus \{i\}$. Therefore $x_i$ is a fan in $\phi'$.

We have thus proved that $\phi'$ is a general position geometric drawing of $G$, such that for each $r \in \{1, 2, \ldots, k\}$, the induced drawing of $G_r$ is noncrossing, and every $2k$-simplicial vertex is a fan. Thus $\phi'$ has thickness $k$ and is a good drawing of $G$. This completes the proof of Proposition 6, which implies Theorem 1.

Note that it is easily seen that each noncrossing subgraph $G_r$ in the proof of Proposition 6 is series-parallel.

9 Book Thickness Lower Bound

Here we prove Theorem 3 for $k \geq 3$. By the discussion in Section 5 it suffices to construct a $k$-tree $Q_k$ with book thickness $bt(Q_k) \geq k + 1$ for all $k \geq 3$. To do so, start with the $k$-tree $K_{k, 2k^2+1}^*$ defined in Section 3. Recall that $K$ is a $k$-clique and $S$ is a set of $2k^2 + 1$ $k$-simplicial vertices in $K_{k, 2k^2+1}^*$. For each vertex $v \in S$, choose three distinct vertices $x_1, x_2, x_3 \in K$, and for each $i \in \{1, 2, 3\}$, add a set of four vertices onto the $k$-clique $(K \cup \{v\}) \setminus \{x_i\}$. Each set of four vertices is called an $i$-block of $v$. Let $T$ be the set of vertices added in this step. Clearly $Q_k$ is a $k$-tree; see Figure 8.

![Figure 8: The graph $Q_k$ in Theorem 3 with $k = 3$.](image)

**Lemma 6.** The book thickness of $Q_k$ satisfies $bt(Q_k) \geq k + 1$. 

23
Proof. Suppose, for the sake of contradiction, that $Q_k$ has a book embedding with thickness $k$. Let $\{E_1, E_2, \ldots, E_k\}$ be the corresponding partition of the edges. For each ordered pair of vertices $v, w \in V(Q_k)$, let the arc-set $V_{vw}$ be the list of vertices in clockwise order from $v$ to $w$ (not including $v$ and $w$).

Say $K = (u_1, u_2, \ldots, u_k)$ in anticlockwise order. There are $2k^2 + 1$ vertices in $S$. Without loss of generality there are at least $2k + 1$ vertices in $S \cap V_{v_1u_k}$. Let $(v_1, v_2, \ldots, v_{2k+1})$ be $2k + 1$ vertices in $S \cap V_{v_1u_k}$ in clockwise order.

Observe that the $k$ edges $\{u_iv_{k+i} : 1 \leq i \leq k\}$ are pairwise crossing, and thus receive distinct colours, as illustrated in Figure 9(a). Without loss of generality, each $u_iv_{k+i} \in E_i$. As illustrated in Figure 9(b), this implies that $u_1v_{2k+1} \in E_1$, since $u_1v_{2k+1}$ crosses all of $\{u_iv_{k+i} : 2 \leq i \leq k\}$ which are coloured $\{2, 3, \ldots, k\}$. As illustrated in Figure 9(c), this in turn implies $u_2v_2 \in E_2$, and so on. By an easy induction, we obtain that $u_iv_{k+2-i} \in E_i$ for all $i \in \{1, 2, \ldots, k\}$, as illustrated in Figure 9(d). It follows that for all $i \in \{1, 2, \ldots, k\}$ and $j \in \{k - i + 1, k - i + 2, \ldots, 2k - 2 - i\}$, the edge $u_iv_j \in E_i$, as illustrated in Figure 9(e). Finally, as illustrated in Figure 9(f), we have:

If $qu_i \in E(Q_k)$ and $q \in V_{v_{k+i}v_{k+i+1}}$, then $qu_i \in E_i$.  

(\ast)

Consider one of the twelve vertices $w \in T$ that are added onto a clique that contain $v_{k+1}$. Then $w$ is adjacent to $v_{k+1}$. Moreover, $w$ is in $V_{v_{k+1}v_{k+1}}$ or $V_{v_{k+1}v_{k+1}}$, as otherwise the edge $wv_{k+1}$ crosses $k$ edges of $Q_k[\{v_{k-1}, v_{k+1}\}; K]$ that are all coloured differently, which is a contradiction. By the pigeon-hole principle, one of $V_{v_{k+1}v_{k+1}}$ and $V_{v_{k+1}v_{k+1}}$ contains at least
two vertices from two distinct $p$-blocks of $v_{k+1}$. Without loss of generality, $V_{v_{k+1}}$ does. Let these four vertices be $(a, b, c, d)$ in clockwise order.

Each vertex in $\{b, c, d\}$ is adjacent to $k - 1$ vertices of $K$. Not all of $b, c, d$ are adjacent to the same subset of $k - 1$ vertices in $K$, as otherwise all of $b, c, d$ would belong to the same $p$-block. Hence each vertex in $K$ has a neighbour in $\{b, c, d\}$. By $(\star)$ the edges of $Q_k[\{b, c, d\}, K]$ receive all $k$ colours. However, every edge in $Q_k[\{b, c, d\}; K]$ crosses the edge $av_{k+1}$, implying that there is no colour available for $av_{k+1}$. This contradiction completes the proof. \hfill \Box

Note that the number of vertices in $Q_k$ is $|K| + |S| + |T| = k + 2k^2 + 1 + 3 \cdot 4 \cdot (2k^2 + 1) = 13(2k^2 + 1) + k$. Adding more simplicial vertices to $Q_k$ cannot reduce its book thickness. Thus for all $n \geq 13(2k^2 + 1) + k$, there is a $k$-tree $G$ with $n$ vertices and $bt(G) = k + 1$.

## 10 Open Problems

### Complete Graphs:

The thickness of the complete graph $K_n$ was intensely studied in the 1960’s and 1970’s. Results by a number of authors [2, 10, 11, 58] together prove that $\theta(K_n) = \lceil (n + 2)/6 \rceil$, unless $n = 9$ or 10, in which case $\theta(K_9) = \theta(K_{10}) = 3$. Bernhart and Kainen [13] proved that $bt(K_n) = \lfloor n/2 \rfloor$. In fact, it is easily seen that

$$a(K_n) = \overline{a}(K_n) = bt(K_n) = ba(K_n) = \lfloor n/2 \rfloor.$$  

Akiyama and Kano [1] proved that $sa(K_n) = \lfloor n/2 \rfloor + 1$. Now

$$\overline{sa}(K_n) \leq bsa(K_n) \leq n - 1.$$  

(Proof. Place the vertices of $K_n$ on a circle, with a spanning star rooted at each vertex except one.) What is $\overline{sa}(K_n)$ and $bsa(K_n)$?

Bose et al. [17] proved that every geometric drawing of $K_n$ has arboricity (and thus thickness) at most $n - \sqrt{n/12}$. It is unknown whether for some constant $\varepsilon > 0$, every geometric drawing of $K_n$ has thickness at most $(1 - \varepsilon)n$; see [17]. Dillencourt et al. [26] studied the geometric thickness of $K_n$, and proved that

$$\lfloor (n/5.646) + 0.342 \rfloor \leq \overline{t}(K_n) \leq \lfloor n/4 \rfloor.$$  

What is $\overline{t}(K_n)$? It seems likely that the answer is closer to $\lfloor n/4 \rfloor$ rather than the above lower bound.

### Asymptotics:

Eppstein [34] (also see [14]) constructed $n$-vertex graphs $G_n$ with $sa(G_n) = a(G_n) = \theta(G_n) = \overline{t}(G_n) = 2$ and $bt(G_n) \to \infty$. Thus book thickness is not bounded by any function of geometric thickness. Similarly, Eppstein [34] constructed $n$-vertex graphs $H_n$ with $sa(H_n) = a(H_n) = \theta(H_n) = 3$ and $\overline{t}(H_n) \to \infty$. Thus geometric thickness is not bounded by any function of thickness (or arboricity). Eppstein [34] asked whether graphs with thickness

---

\footnote{Archdeacon [6] writes, “The question (of the value of $\overline{t}(K_n)$) was apparently first raised by Greenberg in some unpublished work. I read some of his personal notes in the library of the University of Riga in Latvia. He gave a construction that showed $\overline{t}(K_n) \leq \lfloor n/4 \rfloor.$”}
2 have bounded geometric thickness? Whether all graphs with arboricity 2 have bounded geometric thickness is also interesting. It is easily seen that graphs with star arboricity 2 have geometric star arboricity at most 2 (cf. [18]).

**Book Arboricity:** Bernhart and Kainen [13] proved that every graph \( G \) with book thickness \( t \) satisfies \( |E(G)| \leq (t + 1)|V(G)| - 3t \). Thus Equation (1) implies that \( a(G) \leq bt(G) + 1 \) for every graph \( G \), as observed by Dean and Hutchinson [21]. Is \( ba(G) \leq bt(G) + 1 \)?

**Number of Edges:** Let \( E_m \) be the class of graphs with at most \( m \) edges. Dean et al. [22] proved that \( \theta(E_m) \leq \sqrt{m/3} + 3/2 \). What is the minimum \( c \) such that \( \theta(E_m) \leq (c + o(1))\sqrt{m} \)? Dean et al. [22] conjectured that the answer is \( c = 1/16 \), which would be tight for the balanced complete bipartite graph [12]. Malitz [55] proved using a probabilistic argument that \( bt(E_m) \leq 72\sqrt{m} \). Is there a constructive proof that \( bt(E_m) \in O(\sqrt{m}) \) or \( \overline{bt}(E_m) \in O(\sqrt{m}) \)? What is the minimum \( c \) such that \( \overline{bt}(E_m) \leq (c + o(1))\sqrt{m} \) or \( bt(E_m) \leq (c + o(1))\sqrt{m} \)?

**Planar Graphs:** Recall that Yannakakis [78] proved that every planar graph \( G \) has book thickness \( bt(G) \leq 4 \). He also claims there is a planar graph \( G \) with \( bt(G) = 4 \). A construction is given in the conference version of his paper [77], but the proof is far from complete: Yannakakis admits, “Of course, there are many other ways to lay out the graph” [77]. The journal version [78] cites a paper “in preparation” that proves the lower bound. This paper has not been published. Therefore we consider it an open problem whether \( bt(G) \leq 3 \) for every planar graph \( G \).

Let \( G_0 = K_3 \). For \( k \geq 1 \), let \( G_k \) be the planar 3-tree obtained by adding a 3-simplicial vertex onto the vertex set of each face of \( G_{k-1} \). In the journal version of this paper we conjectured that \( bt(G_k) = 4 \) for sufficiently large \( k \). This conjecture is false. Indeed, Heath and Istrail [45], Malitz [54] proved using a probabilistic argument that \( bt(S) \in O(\sqrt{\gamma}) \), and thus \( \overline{bt}(S) \in O(\sqrt{\gamma}) \). Is there a constructive proof that \( bt(S) \in O(\sqrt{\gamma}) \) or \( \overline{bt}(S) \in O(\sqrt{\gamma}) \)? What is the minimum \( c \) such that \( bt(S) \leq (c + o(1))\sqrt{\gamma} \), or \( \overline{bt}(S) \leq (c + o(1))\sqrt{\gamma} \)?

Endo [32] proved that \( bt(S_1) \leq 7 \). Let \( \chi(S_\gamma) \) denote the maximum chromatic number of all graphs with genus at most \( \gamma \). Dean and Hutchinson [21] proved that \( \theta(S_\gamma) \leq 6 + 2\sqrt{\gamma - 2} \); also see [2, 8]. What is the minimum \( c \) such that \( \theta(S_\gamma) \leq (c + o(1))\sqrt{\gamma} \)? Building on prior work by Heath and Istrail [45], Malitz [54] proved using a probabilistic argument that \( bt(S_\gamma) \in O(\sqrt{\gamma}) \), and thus \( \overline{bt}(S_\gamma) \in O(\sqrt{\gamma}) \). Is there a constructive proof that \( bt(S_\gamma) \in O(\sqrt{\gamma}) \) or \( \overline{bt}(S_\gamma) \in O(\sqrt{\gamma}) \)? What is the minimum \( c \) such that \( bt(S_\gamma) \leq (c + o(1))\sqrt{\gamma} \), or \( \overline{bt}(S_\gamma) \leq (c + o(1))\sqrt{\gamma} \)?

Endo [32] asked whether \( bt(S_\gamma) = \chi(S_\gamma) \) for all \( \gamma \). Both \( bt(S_\gamma) \) and \( \chi(S_\gamma) \) are in \( O(\sqrt{\gamma}) \). There is some tangible evidence relating book thickness and chromatic number. First, Bernhart and Kainen [13] proved that \( \chi(G) \leq 2 \cdot bt(G) + 2 \) for every graph \( G \). Second, the maximum book thickness and maximum chromatic number coincide \((= k + 1)\) for
graphs of treewidth \( k \geq 3 \). In fact, the proof by Ganley and Heath [37] that \( \text{bt}(T_k) \leq k + 1 \) is based on the \((k + 1)\)-colourability of \( k \)-trees.

Minors: Let \( M_\ell \) be the class of graphs with no \( K_\ell \)-minor. Note that \( M_3 = \mathcal{T}_1 \) and \( M_4 = \mathcal{T}_2 \). Jünger et al. [49] proved that \( \theta(M_5) = 2 \). What is \( \overline{\theta}(M_5) \) and \( \text{bt}(M_5) \)? Kostochka [53] and Thomason [70] independently proved that the maximum arboricity of all graphs with no \( K_\ell \)-minor is \( \Theta(\ell \sqrt{\log \ell}) \). In fact, Thomason [71] asymptotically determined the right constant. Thus \( \theta(M_\ell) \in \Theta(\ell \sqrt{\log \ell}) \) by Equation (2). Blankenship and Oporowski [14, 15] proved that \( \text{bt}(M_\ell) \) (and hence \( \overline{\theta}(M_\ell) \)) is finite. The proof depends on Robertson and Seymour’s deep structural characterisation of the graphs in \( M_\ell \). As a result, the bound on \( \text{bt}(M_\ell) \) is a truly huge function of \( \ell \). Is there a simple proof that \( \overline{\theta}(M_\ell) \) or \( \text{bt}(M_\ell) \) is finite? What is the right order of magnitude of \( \overline{\theta}(M_\ell) \) and \( \text{bt}(M_\ell) \)?

Maximum Degree: Let \( D_\Delta \) be the class of graphs with maximum degree at most \( \Delta \). Wessel [75] and Halton [44] independently proved that \( \theta(D_\Delta) \leq \lceil \Delta/2 \rceil \), and Sýkora et al. [68] proved that \( \theta(D_\Delta) \geq \lceil \Delta/2 \rceil \). Thus \( \theta(D_\Delta) = \lceil \Delta/2 \rceil \). Eppstein [34] asked whether \( \overline{\theta}(D_\Delta) \) is finite. A positive result in this direction was obtained by Duncan et al. [29], who proved that \( \overline{\theta}(D_4) \leq 2 \). On the other hand, Barát et al. [9] recently proved that \( \overline{\theta}(D_\Delta) = \infty \) for all \( \Delta \geq 9 \); in particular, there exists \( \Delta \)-regular \( n \)-vertex graphs with geometric thickness \( \Omega(\sqrt{\Delta n^{1/2 - \epsilon}}) \). It is unknown whether \( \overline{\theta}(D_\Delta) \) is finite for \( \Delta \in \{5, 6, 7, 8\} \).

Malitz [55] proved that there exists \( \Delta \)-regular \( n \)-vertex graphs with book thickness \( \Omega(\sqrt{\Delta n^{1/2 - 1/2}}) \). Barát et al. [9] reached the same conclusion for all \( \Delta \geq 3 \). Thus \( \text{bt}(D_\Delta) = \infty \) unless \( \Delta \leq 2 \). Open problems remain for specific values of \( \Delta \). For example, the best bounds on \( \text{bt}(D_3) \) are \( \Omega(n^{1/6}) \) and \( O(n^{1/2}) \).

Computational Complexity: Arboricity can be computed in polynomial time using the matroid partitioning algorithm of Edmonds [30]. Computing the thickness of a graph is \( \mathcal{NP} \)-hard [56]. Testing whether a graph has book thickness at most 2 is \( \mathcal{NP} \)-complete [76]. Dillencourt et al. [26] asked what is the complexity of determining the geometric thickness of a given graph? The same question can be asked for all of the other parameters discussed in this paper.

Acknowledgement

Thanks to the anonymous reviewer for numerous helpful comments.

References

[1] Jin Akiyama and Mikio Kano. Path factors of a graph. In Graphs and applications, pp. 1–21. Wiley, 1985.

[2] V. B. Alekseev and V. S. Gonchakov. Thickness of arbitrary complete graphs. Mat. Sbornik, 101:212–230, 1976.
[3] I. ALGOR AND NOGA ALON. The star arboricity of graphs. *Discrete Math.*, 75(1-3):11–22, 1989.

[4] NOGA ALON, COLIN McDiARMID, AND BRUCE REED. Star arboricity. *Combinatorica*, 12(4):375–380, 1992.

[5] YASUKAZU AOKI. The star-arboricity of the complete regular multipartite graphs. *Discrete Math.*, 81(2):115–122, 1990.

[6] DAN ARCHDEACON. The geometric thickness of the complete graph, 2003. [http://www.emba.uvm.edu/~archdeac/problems/geomthick.html](http://www.emba.uvm.edu/~archdeac/problems/geomthick.html)

[7] KOUHEI ASANO. On the genus and thickness of graphs. *J. Combin. Theory Ser. B*, 43(3):287–292, 1987.

[8] KOUHEI ASANO. On the thickness of graphs with genus 2. *Ars Combin.*, 38:87–95, 1994.

[9] JÁNOS BARÁT, JIŘÍ MATOUŠEK, AND DAVID R. WOOD. Bounded-degree graphs have arbitrarily large geometric thickness. *Electron. J. Combin.*, 13(1):R3, 2006.

[10] LOWELL W. BEINEKE. The decomposition of complete graphs into planar subgraphs. In FRANK HARARY, ed., *Graph Theory and Theoretical Physics*, pp. 139–154. Academic Press, 1967.

[11] LOWELL W. BEINEKE AND FRANK HARARY. The thickness of the complete graph. *Canad. J. Math.*, 17:850–859, 1965.

[12] LOWELL W. BEINEKE, FRANK HARARY, AND J. W. MOON. On the thickness of the complete bipartite graph. *Proc. Cambridge Philos. Soc.*, 60(1):1–5, 1964.

[13] FRANK R. BERNHART AND PAUL C. KAINE. The book thickness of a graph. *J. Combin. Theory Ser. B*, 27(3):320–331, 1979.

[14] ROBIN BLANKENSHIP. *Book Embeddings of Graphs*. Ph.D. thesis, Department of Mathematics, Louisiana State University, U.S.A., 2003.

[15] ROBIN BLANKENSHIP AND BOGDAN OPOROWSKI. Book embeddings of graphs and minor-closed classes. In *Proc. 32nd Southeastern International Conf. on Combinatorics, Graph Theory and Computing*. Department of Mathematics, Louisiana State University, 2001.

[16] HANS L. BODLAENDER. A partial k-arboretum of graphs with bounded treewidth. *Theoret. Comput. Sci.*, 209(1-2):1–45, 1998.

[17] PROSENJIT BOSE, FERRAN HURTADO, EDUARDO RIVERA-CAMPO, AND DAVID R. WOOD. Partitions of complete geometric graphs into plane trees. *Comput. Geom.*, 34(2):116–125, 2006.
[18] Peter Brass, Eowyn Čenek, Christian A. Duncan, Alon Efrat, Cesim Erten, Dan Ismailescu, Stephen G. Kobourov, Anna Lubiw, and Joseph S. B. Mitchell. On simultaneous planar graph embeddings. *Comput. Geom.*, 36(2):117–130, 2007.

[19] Peter Brass, William O. J. Moser, and János Pach. *Research problems in discrete geometry*. Springer, New York, 2005. ISBN 0-387-23815-8.

[20] Dilip Chhajed. Edge coloring a $k$-tree into two smaller trees. *Networks*, 29(4):191–194, 1997.

[21] Alice M. Dean and Joan P. Hutchinson. Relations among embedding parameters for graphs. In *Graph theory, combinatorics, and applications, Vol. 1*, pp. 287–296. Wiley, 1991.

[22] Alice M. Dean, Joan P. Hutchinson, and Edward R. Scheinerman. On the thickness and arboricity of a graph. *J. Combin. Theory Ser. B*, 52(1):147–151, 1991.

[23] Giuseppe Di Battista, Peter Eades, Roberto Tamassia, and Ioannis G. Tollis. *Graph Drawing: Algorithms for the Visualization of Graphs*. Prentice-Hall, 1999. ISBN 0-13-301615-3.

[24] Emilio Di Giacomo, Walter Didimo, Giuseppe Liotta, and Stephen K. Wismath. Book embeddability of series-parallel digraphs. *Algorithmica*, 45(4):531–547, 2006.

[25] Reinhard Diestel. *Graph theory*, vol. 173 of *Graduate Texts in Mathematics*. Springer, 2nd edn., 2000. ISBN 0-387-95014-1.

[26] Michael B. Dillencourt, David Eppstein, and Daniel S. Hirschberg. Geometric thickness of complete graphs. *J. Graph Algorithms Appl.*, 4(3):5–17, 2000.

[27] Guoli Ding, Bogdan Oporowski, Daniel P. Sanders, and Dirk Vertigan. Partitioning graphs of bounded tree-width. *Combinatorica*, 18(1):1–12, 1998.

[28] Vida Dujmović and David R. Wood. On linear layouts of graphs. *Discrete Math. Theor. Comput. Sci.*, 6(2):339–358, 2004.

[29] Christian A. Duncan, David Eppstein, and Stephen G. Kobourov. The geometric thickness of low degree graphs. In *Proc. 20th ACM Symp. on Computational Geometry (SoCG ’04)*, pp. 340–346. ACM Press, 2004.

[30] Jack Edmonds. Minimum partition of a matroid into independent subsets. *J. Res. Nat. Bur. Standards Sect. B*, 69B:67–72, 1965.

[31] Ehab S. El-Mallah and Charles J. Colbourn. Partitioning the edges of a planar graph into two partial $k$-trees. In *Proc. 19th Southeastern Conf. on Combinatorics, Graph Theory, and Computing*, vol. 66 of *Congr. Numer.*., pp. 69–80. 1988.
[32] Toshiki Endo. The pagewidth of toroidal graphs is at most seven. *Discrete Math.*, 175(1-3):87–96, 1997.

[33] David Eppstein. Separating geometric thickness from book thickness, 2001. [http://arXiv.org/math/0109195](http://arXiv.org/math/0109195).

[34] David Eppstein. Separating thickness from geometric thickness. In János Pach, ed., *Towards a Theory of Geometric Graphs*, vol. 342 of *Contemporary Mathematics*, pp. 75–86. Amer. Math. Soc., 2004.

[35] István Fáry. On straight line representation of planar graphs. *Acta Univ. Szeged. Sect. Sci. Math.*, 11:229–233, 1948.

[36] Guillaume Fertin, André Raspaud, and Bruce Reed. On star coloring of graphs. *J. Graph Theory*, 47(3):163–182, 2004.

[37] Joseph L. Ganley and Lenwood S. Heath. The pagewidth of k-trees is O(k). *Discrete Appl. Math.*, 109(3):215–221, 2001.

[38] Daniel Gonçalves. Edge partition of planar graphs into two outerplanar graphs. In *Proc. 37th ACM Symp. on Theory of Computing* (STOC ’05), pp. 504–512. ACM, 2005.

[39] Barry Guiduli. On incidence coloring and star arboricity of graphs. *Discrete Math.*, 163(1-3):275–278, 1997.

[40] Richard K. Guy. Outerthickness and outercoarseness of graphs. In *Proc. British Combinatorial Conf.*, vol. 13 of *London Math. Soc. Lecture Note Ser.*, pp. 57–60. Cambridge Univ. Press, 1974.

[41] Richard K. Guy and Richard J. Nowakowski. The outerthickness & outercoarseness of graphs. I. The complete graph & the n-cube. In *Topics in combinatorics and graph theory*, pp. 297–310. Physica, Heidelberg, 1990.

[42] Richard K. Guy and Richard J. Nowakowski. The outerthickness & outercoarseness of graphs. II. The complete bipartite graph. In *Contemporary methods in graph theory*, pp. 313–322. Bibliographisches Inst., Mannheim, 1990.

[43] S. Louis Hakimi, John Mitchem, and Edward Schmeichel. Star arboricity of graphs. *Discrete Math.*, 149(1-3):93–98, 1996.

[44] John H. Halton. On the thickness of graphs of given degree. *Inform. Sci.*, 54(3):219–238, 1991.

[45] Lenwood S. Heath and Sorin Istrail. The pagewidth of genus g graphs is O(g). *J. Assoc. Comput. Mach.*, 39(3):479–501, 1992.

[46] Arthur M. Hobbs. A survey of thickness. In *Recent Progress in Combinatorics* (Proc. 3rd Waterloo Conf. on Combinatorics, 1968), pp. 255–264. Academic Press, New York, 1969.
[47] Wen-Ting Huang. *Linear Arboricity and Star Arboricity of Graphs*. Ph.D. thesis, National Central University, Taiwan, 2003.

[48] Joan P. Hutchinson, Thomas C. Shermer, and Andrew Vince. On representations of some thickness-two graphs. *Comput. Geom.*, 13(3):161–171, 1999.

[49] Michael Jünger, Petra Mutzel, Thomas Odenthal, and Mark Scharbrodt. The thickness of a minor-excluded class of graphs. *Discrete Math.*, 182(1-3):169–176, 1998.

[50] Paul C. Kainen. Thickness and coarseness of graphs. *Abh. Math. Sem. Univ. Hamburg*, 39:88–95, 1973.

[51] Michael Kaufmann and Dorothea Wagner, eds. *Drawing Graphs: Methods and Models*, vol. 2025 of Lecture Notes in Comput. Sci. Springer, 2001.

[52] Kiran S. Kedlaya. Outerplanar partitions of planar graphs. *J. Combin. Theory Ser. B*, 67(2):238–248, 1996.

[53] Alexandr V. Kostochka. The minimum Hadwiger number for graphs with a given mean degree of vertices. *Metody Diskret. Analiz.*, 38:37–58, 1982.

[54] Seth M. Malitz. Genus g graphs have pagенumber $O(\sqrt{g})$. *J. Algorithms*, 17(1):85–109, 1994.

[55] Seth M. Malitz. Graphs with $E$ edges have pagенumber $O(\sqrt{E})$. *J. Algorithms*, 17(1):71–84, 1994.

[56] Anthony Mansfield. Determining the thickness of graphs is NP-hard. *Math. Proc. Cambridge Philos. Soc.*, 93(1):9–23, 1983.

[57] Jiří Matoušek. *Lectures on Discrete Geometry*, vol. 212 of Graduate Texts in Mathematics. Springer, 2002. ISBN 0-387-95373-6.

[58] Jean Mayer. Décomposition de $K_{16}$ en trois graphes planaires. *J. Combinatorial Theory Ser. B*, 13:71, 1972.

[59] James R. Munkres. *Topology: a first course*. Prentice-Hall, 1975. ISBN 0-139-25495-1.

[60] Petra Mutzel, Thomas Odenthal, and Mark Scharbrodt. The thickness of graphs: a survey. *Graphs Combin.*, 14(1):59–73, 1998.

[61] Crispin St. J. A. Nash-Williams. Decomposition of finite graphs into forests. *J. London Math. Soc.*, 39:12, 1964.

[62] Takao Nishizeki and Md. Saidur Rahman. *Planar graph drawing*, vol. 12 of Lecture Notes Series on Computing. World Scientific, 2004. ISBN 981-256-033-5.
[63] L. Taylor Ollmann. On the book thicknesses of various graphs. In Frederick Hoffman, Roy B. Levow, and Robert S. D. Thomas, eds., Proc. 4th Southeastern Conference on Combinatorics, Graph Theory and Computing, vol. VIII of Congressus Numerantium, p. 459. 1973.

[64] János Pach and Rephael Wenger. Embedding planar graphs at fixed vertex locations. Graphs Combin., 17(4):717–728, 2001.

[65] Timo Poranen and Erkki Mäkinen. Remarks on the thickness and outerthickness of a graph. Comput. Math. Appl., 50(1-2):249–254, 2005.

[66] Bruce A. Reed. Algorithmic aspects of tree width. In Bruce A. Reed and Cláudia L. Sales, eds., Recent Advances in Algorithms and Combinatorics, pp. 85–107. Springer, 2003.

[67] S. Rengarajan and C. E. Veni Madhavan. Stack and queue number of 2-trees. In Ding-Zhu Du and Ming Li, eds., Proc. 1st Annual International Conf. on Computing and Combinatorics (COCOON ’95), vol. 959 of Lecture Notes in Comput. Sci., pp. 203–212. Springer, 1995.

[68] Ondrej Sýkora, László A. Székely, and Imrich Vrťo. A note on Halton’s conjecture. Inform. Sci., 164(1-4):61–64, 2004.

[69] László A. Székely. A successful concept for measuring non-planarity of graphs: the crossing number. Discrete Math., 276(1-3):331–352, 2004.

[70] Andrew Thomason. An extremal function for contractions of graphs. Math. Proc. Cambridge Philos. Soc., 95(2):261–265, 1984.

[71] Andrew Thomason. The extremal function for complete minors. J. Combin. Theory Ser. B, 81(2):318–338, 2001.

[72] Mitsunori Togasaki and Koichi Yamazaki. Pagenumber of pathwidth-k graphs and strong pathwidth-k graphs. Discrete Math., 259(1-3):361–368, 2002.

[73] William T. Tutte. The thickness of a graph. Nederl. Akad. Wetensch. Proc. Ser. A 66=Indag. Math., 25:567–577, 1963.

[74] Klaus Wagner. Bemerkung zum Vierfarbenproblem. Jber. Deutsch. Math.-Verein., 46:26–32, 1936.

[75] Walter Wessel. Über die Abhängigkeit der Dicke eines Graphen von seinen Knotenpunkten. In 2nd Colloquium for Geometry and Combinatorics, pp. 235–238. Tech. Hochschule Karl-Marx-Stadt, 1983.

[76] Avi Wigderson. The complexity of the Hamiltonian circuit problem for maximal planar graphs. Tech. Rep. EECS 198, Princeton University, U.S.A., 1982.
Appendix

Here we prove two perturbation lemmas that are used in the proof of Theorem 1.

Lemma 7. Let $P$ be a finite set of points in the plane. Then there exists $\varepsilon > 0$ such that
(a) $D_\varepsilon(u) \cap D_\varepsilon(v) = \emptyset$ for all $u, v \in P$,
(b) for all $u, v, w \in P$, if the discs $D_\varepsilon(u), D_\varepsilon(v), D_\varepsilon(w)$ are collinear, then the points $u, v, w$ are collinear.

Proof. Say $P = \{p_1, \ldots, p_n\}$. We prove the following statement by induction on $\ell \in \{0, 1, \ldots, n\}$:

For all $i \in \{1, 2, \ldots, \ell\}$, there exists a disk $D^\ell(p_i)$ of positive radius centered at $p_i$ such that the following two properties hold, where $D^\ell(p_i) := \{p_i\}$ for each $i > \ell$:
(a) $D^\ell(p_i) \cap D^\ell(p_j) = \emptyset$ for all $p_i, p_j \in P$,
(b) for all $p_i, p_j, p_k \in P$, if the discs $D^\ell(p_i), D^\ell(p_j), D^\ell(p_k)$ are collinear, then the points $p_i, p_j, p_k$ are collinear.

This statement implies the lemma, by defining $\varepsilon$ to be the radius of the smallest disk $D^0(p_i)$.

The base case $\ell = 0$ is vacuous. Now assume that $\ell > 0$. For all distinct $i$ and $j$ such that $p_i, p_j, p_k$ are not collinear, every line that intersects $D^{\ell-1}(p_i)$ and $D^{\ell-1}(p_j)$ does not intersect $p_k$, by induction. Thus $p_k$ is in an open region $R$ of the plane defined by the complement of the union of all such lines. Thus, there is an open disk $D \subset R$ of positive radius centered at $p_k$, such that $D$ does not intersect $D^{\ell-1}(p_i)$, for all $i \neq \ell$. Defining $D^\ell(p_k) := D$ and $D^\ell(p_i) := D^{\ell-1}(p_i)$ for all points $p_i \neq p_k$ completes the proof.

Lemma 8. Let $\phi$ be a noncrossing geometric drawing of a graph $G$ (not necessarily in general position). Then there exists $\varepsilon > 0$ such that if $\phi' : V(G) \to \mathbb{R}^2$ is an injection with $\phi'(v) \in D_\varepsilon(\phi(v))$ for every vertex $v \in V(G)$, then $\phi'$ is a noncrossing geometric drawing of $G$ with the property that if three vertices $\phi'(u), \phi'(v), \phi'(w)$ are collinear in $\phi'$, then $\phi(u), \phi(v), \phi(w)$ are collinear in $\phi$.

Proof. Let $\varepsilon > 0$ be the constant obtained by applying Lemma 7 to the point set $\{\phi(v) : v \in V(G)\}$. We now prove that every function $\phi'$, as defined in the statement of the lemma, has the desired properties.

Suppose there exists a line intersecting three vertices $\phi'(u), \phi'(v), \phi'(w)$ in $\phi'$. Then the same line passes through the discs $D_\varepsilon(\phi(u)), D_\varepsilon(\phi(v)), D_\varepsilon(\phi(w))$. Thus, by Lemma 7, $\phi(u), \phi(v), \phi(w)$ are collinear in $\phi$.

It remains to prove that $\phi'$ is a noncrossing geometric drawing of $G$. For each vertex $v \in V(G)$, let $J(v) := D_\varepsilon(\phi(v))$. Thus the image $\phi'(v)$ is in $J(v)$. For each edge $vw \in E(G)$. 

[77] Mihalis Yannakakis. Four pages are necessary and sufficient. In Proc. 18th ACM Symp. on Theory of Comput. (STOC ’86), pp. 104–108. 1986.

[78] Mihalis Yannakakis. Embedding planar graphs in four pages. J. Comput. System Sci., 38(1):36–67, 1989.
let $J(vw)$ be the region consisting of the union of all segments with one endpoint in $J(v)$ and the other endpoint in $J(w)$. Since $\phi'(v) \in J(v)$ and $\phi'(w) \in J(w)$, the image $\overline{\phi'(v)\phi'(w)}$ of the edge $vw$ is contained in $J(vw)$.

Thus to prove that $\phi'$ is a noncrossing geometric drawing of $G$, it suffices to prove that:

(i) $J(v) \cap J(xy) = \emptyset$, for every vertex $v \in V(G)$ and edge $xy \in E(G)$ not incident to $v$, and

(ii) $J(vw) \cap J(xy) = \emptyset$, for all edges $vw, xy \in E(G)$ with no common endpoint.

We now prove (i). First, suppose that $\phi(v), \phi(x), \phi(y)$ are collinear. Then without loss of generality, $\phi(x)$ is between $\phi(v)$ and $\phi(y)$, as otherwise $v$ intercepts the edge $xy$ in $\phi$. Since $J(v) \cap J(x) = \emptyset$ by Lemma 7, we have $J(v) \cap J(xy) = \emptyset$. Now suppose that $\phi(v), \phi(x), \phi(y)$ are not collinear. By Lemma 7, $J(v), J(x), J(y)$ are not collinear, in which case $J(v) \cap J(xy) = \emptyset$.

We now prove (ii). Suppose on the contrary that $J(vw) \cap J(xy) \neq \emptyset$. First suppose that no three points in $\{\phi(v), \phi(w), \phi(x), \phi(y)\}$ are collinear. Then at least one of $vw$ and $xy$, say $vw$, lies on the convex hull of $\{\phi(v), \phi(w), \phi(x), \phi(y)\}$. Then $\overline{\phi(v)\phi(w)}$ does not intersect $\overline{\phi(x)\phi(y)}$. Therefore, the only way for $J(vw) \cap J(xy) \neq \emptyset$ is if $J(x) \cup J(v) \cap J(vw) \neq \emptyset$ or $J(v) \cup J(w) \cap J(xy) \neq \emptyset$. This is impossible by Lemma 7 since no three points in $\{\phi(v), \phi(w), \phi(x), \phi(y)\}$ are collinear.

Now suppose that exactly three vertices in $\{\phi(v), \phi(w), \phi(x), \phi(y)\}$ are collinear. Without loss of generality, $\phi(v), \phi(w), \phi(x)$ are collinear and $\phi(w)$ is between $\phi(v)$ and $\phi(x)$. Then the only way for $J(vw) \cap J(xy) \neq \emptyset$ is if $J(w) \cap J(xy) \neq \emptyset$, which is impossible by Lemma 7 since $\phi(w), \phi(x), \phi(y)$ are not collinear by assumption.

Finally assume that all four points $\phi(v), \phi(w), \phi(x), \phi(y)$ are collinear. Since $vw$ and $xy$ do not cross in $\phi$, we may assume without loss of generality, that $v, w, x, y$ are in this order on the line. Then $J(vw) \cap J(xy) \neq \emptyset$ only if $J(w) \cap J(x) \neq \emptyset$, which is impossible by Lemma 7.  

34