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On the CVP for the root lattices via folding with deep ReLU neural networks

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Abstract—Point lattices and their decoding via neural networks are considered in this paper. Lattice decoding in \(\mathbb{R}^n\), known as the closest vector problem (CVP), becomes a classification problem in the fundamental parallelotope with a piecewise linear function defining the boundary. Theoretical results are obtained by studying root lattices. We show how the number of pieces in the boundary function reduces dramatically with folding, from exponential to linear. This translates into a two-layer ReLU network requiring a number of neurons growing exponentially in \(n\) to solve the CVP, whereas this complexity becomes polynomial in \(n\) for a deep ReLU network.

I. INTRODUCTION AND MOTIVATIONS

The objective of this paper is two-fold. Firstly, we introduce a new paradigm to solve the CVP. This approach enables to find efficient decoding algorithms for some dense lattices. For instance, such a neural network for the Gosset lattice is a key component of a neural Leech decoder. Secondly, we also aim at contributing to the understanding of the efficiency of deep learning, namely the expressive power of deep neural networks. As a result, our goal is to present new decoding algorithms and interesting functions that can be efficiently computed by deep neural networks.

Deep Learning is about two key aspects: (i) Finding a function class \(\Phi = \{f\}\) that contains a function \(f^*\) “close enough” to a target function. (ii) Finding a learning algorithm \(L\) for the class \(\Phi\). Of course the choices of (i) and (ii) can be either done jointly or separately but in either case they impact each other. Research on the expressive power of deep neural networks focuses mostly on (i) [7][11], by studying some specific functions contained in the function class of a network. Typically, the aim is to show that there exist functions that can be well approximated by a deep network with a polynomial number of parameters whereas an exponential number of parameters is required for a shallow network. This line of work leads to “gap” theorems and “capacity” theorems for deep networks and it is similar to the classical theory of Boolean circuit complexity [8].

In this scope, several papers investigate specifically deep ReLU networks [10][9][11][14][12][2] (See [9, Section 2.1] for a short introduction to ReLU neural networks). Since a ReLU network computes a composition of piecewise affine functions, all functions in \(\Phi\) are continuous piecewise linear (CPWL). Hence, the efficiency of \(\Phi\) can be evaluated by checking whether a CPWL function with a lot of affine pieces belongs to \(\Phi\). For example, there exists at least one function in \(\mathbb{R}^n\) with \(\Omega \left( \frac{(w/n)^{L-1} w^n}{n} \right)\) affine pieces that can be computed with a \(w\)-wide deep ReLU network having \(L\) hidden layers [9]. A two-layer network would need an exponential number of parameters for this same function. In [11] they further show that any random deep network achieves a similar exponential behavior.

Some results in the literature are established by considering elementary oscillatory functions (see e.g. [15] - VII-F) or piecewise linear functions with regions of random shape. It is not clear whether such types of functions may arise naturally in computer science and engineering fields.

Our work lies somewhere between [9][14] and [11]: our functions are neither elementary nor random. We discovered them in the context of sphere packing and lattices which are solution to many fundamental problems in number theory, chemistry, communication theory, string theory, and cryptography [4]. Hence, in contrary to existing works, we do not search for a specific class of functions to justify the success of deep learning. We set technical problems, study the functions arising from lattices, and show that deep networks are suited to tackle them as we obtain similar gap theorems between shallow and deep models.

II. MAIN RESULTS

A ReLU network with a finite number of neurons is not capable to infer infinite periodic functions (see [15] - VII-F). Hence, it cannot implement a simple modulo operation. As a result, we allow a low complexity pre-processing of the point to be decoded to obtain an equivalent position in the fundamental parallelotope \(\mathcal{P}(\mathcal{B})\) of the lattice.

1) Theorems 1&2 show that the decision boundary for the hyperplane logical decoder (HLD) [5] is given by a continuous piecewise linear function, for any lattice with a Voronoi-reduced basis. Corollary 1 guarantees the same result for a semi-Voronoi-reduced basis.
2) For the lattice \(A_n\) with a basis defined by the Gram matrix (4) and a point in \(\mathcal{P}(\mathcal{B})\), Theorem 4 proves that the decision function has \(\Omega(2^n)\) affine pieces.
3) Also for \(A_n\), the number of pieces is reduced to \(\mathcal{O}(n)\) after folding as stated in Theorem 5.
4) Results of Section V-B and V-C, based on Theorem 5, implies that there exists a ReLU network of depth \(O(n^2)\) and width \(O(n^2)\) solving the CVP.
5) Theorem 6 shows that a ReLU network with only one hidden layer needs \(\Omega(2^n)\) neurons to solve the CVP.

Moreover, the theory presented in this paper is not limited to \(A_n\). It extends very well to other dense lattices. Indeed, we
already obtained similar results for all root lattices. They will however be presented in future communications due to lack of space. Finally, this paradigm seems not to be limited to lattices as it may extend to binary block codes (see e.g. Figure 4 in [15] - VII-F).

III. LATTICES AND POLYTOPES

A lattice \( \Lambda \) is a discrete additive subgroup of \( \mathbb{R}^n \). For a rank-\( n \) lattice in \( \mathbb{R}^n \), the rows of a \( n \times n \) generator matrix \( G \) constitute a basis of \( \Lambda \) and any lattice point \( x \) is obtained via \( x = zG \), where \( z \in \mathbb{Z}^n \). For a given basis \( B = \{ b_i \}_{i=1}^n \), \( \mathcal{P}(B) \) denotes the fundamental parallelotope of \( \Lambda \) and \( \mathcal{V}(x) \) the Voronoi cell of a lattice point \( x \) [5]. The minimum Euclidean distance of \( \Lambda \) is \( d_{\min}(\Lambda) = 2\rho \), where \( \rho \) is the packing radius.

A vector \( v \in \Lambda \) is called Voronoi vector if the half-space \( \{ y \in \mathbb{R}^n : \ y \cdot v \leq \frac{1}{2} v \cdot v \} \) has a non empty intersection with \( \mathcal{V}(0) \). This vector is said relevant if the intersection is an \( n \)-1-dimensional face of \( \mathcal{V}(0) \). We denote by \( \tau v \) the number of relevant Voronoi vectors, referred to as the Voronoi number in the sequel. For root lattices [4], the Voronoi number is equal to the kissing number \( \tau \). For random lattices, we typically have \( \tau v = 2^n+1 \). The set \( \tau v(x) \), for \( x \in \Lambda \), is the set of lattice points having a common Voronoi facet with \( x \).

The next definition, introduced in [5], is important for the rest of the paper.

Definition 1. Let \( B \) be the \( \mathbb{Z} \)-basis of a rank-\( n \) lattice \( \Lambda \) in \( \mathbb{R}^n \). \( B \) is said Voronoi-reduced (VR) if, for any point \( y \in \mathcal{P}(B) \), the closest lattice point \( \hat{x} \) to \( y \) is one of the \( 2^n \) corners of \( \mathcal{P}(B) \), i.e. \( \hat{x} = zG \) where \( z \in \{ 0, 1 \}^n \).

Lattice decoding refers to the method of finding the closest lattice point, the closest in Euclidean distance sense. This problem is also known as the closest vector problem. The neural lattice decoder employs \( \mathcal{P}(B) \) as its main compact region [5], thus it is important to characterize \( \mathcal{P}(B) \) as made below.

A \( k \)-dimensional element of \( \mathcal{P}(B) \) is referred to as \( k \)-face. \( \mathcal{P}(B) \) has \( 2^n \) 0-faces, called corners or vertices. This set of corners is denoted \( C_{\mathcal{P}(B)} \). Moreover, the subset of \( C_{\mathcal{P}(B)} \) obtained with \( z_i = 1 \) is \( C_{\mathcal{P}(B)}^{1} \) and \( C_{\mathcal{P}(B)}^{0} \) for \( z_i = 0 \). The remaining faces of \( \mathcal{P}(B) \) are parallelotopes. For instance, a \( n \)-1-dimensional facet of \( \mathcal{P}(B) \), say \( F_i \), is itself a parallelotope of dimension \( n-1 \) defined by \( n-1 \) vectors of \( B \). Throughout the paper, the term facet refers to a \( n \)-1-face.

A polyhedron is defined as the intersection of a finite number of half-spaces generated by hyperplanes:

\[
P_o = \{ x \in \mathbb{R}^n : \ ax \leq b, \ A \in \mathbb{R}^{n \times m}, \ b \in \mathbb{R}^m \}.
\]

A polyhedron is convex if none of the bounding hyperplanes crosses its interior. A convex polyhedron is called a polytope [6]. In this paper, we use not only parallelotopes but also simplices. A \( n \)-simplex associated to \( B \) is given by

\[
S(B) = \{ y \in \mathbb{R}^n : \ y = \sum_{i=0}^n \alpha_i b_i, \sum_{i=1}^n \alpha_i \leq 1, \ \alpha_i \geq 0 \ \forall \ i \}.
\]

It is clear that the corners of \( S(B) \), the set \( C_{S(B)} \), are the \( n+1 \) points \( \{ 0, b_1, ..., b_n \} \).

We say that a function \( f : \mathbb{R}^{n-1} \to \mathbb{R} \) is continuous piece-wise linear (CPWL) if there exists a finite set of polyhedra covering \( \mathbb{R}^{n-1} \) (which implies continuity), and \( g \) is affine over each polyhedron. The number of pieces of \( g \) is the number of distinct expressions of the local affine functions.

Finally, \( \vee \) and \( \wedge \) denote respectively the maximum and the minimum operator. We define a convex (resp. concave) CPWL function formed by a set of affine functions related by the operator \( \vee \) (resp. \( \wedge \)). If \( \{ g_k \} \) is a set of \( K \) affine functions, the function \( f = g_1 \vee ... \vee g_K \) is CPWL and convex.

IV. THE DECISION BOUNDARY FUNCTION

Given a VR basis, after translating the point to be decoded inside \( \mathcal{P}(B) \), the HLD decoder proceeds in estimating each \( z_i \)-component separately [5]. The HLD computes the position of \( \mathcal{P}(B) \) relative to a boundary via a Boolean equation to guess whether \( z_i = 0 \), i.e. the closest lattice point belongs to \( C_{\mathcal{P}(B)}^{1} \), or \( z_i = 1 \) when the closest lattice point is in \( C_{\mathcal{P}(B)}^{0} \). This boundary cuts \( \mathcal{P}(B) \) into two regions. It is composed of Voronoi facets of the corner points. The next step is to study the decision boundary function. For the rest of the paper, without loss of generality, the integer coordinate to be decoded is \( z_1 \).

We recall that a variable \( u_j(y) \) in the Boolean equations of the HLD is obtained as:

\[
u_j(y) = \text{sign}(y \cdot v_j - p_j) \in \{0, 1\}, \quad (1)
\]

where \( v_j \) is the orthogonal vector to the boundary hyperplane \( \{ y \in \mathbb{R}^n : \ y \cdot v_j - p_j = 0 \} \). The latter contains the Voronoi facet of a point \( x \in C_{\mathcal{P}(B)}^{1} \) and a point from \( \tau v_j(x) \cap C_{\mathcal{P}(B)}^{0} \). The decision boundary cutting \( \mathcal{P}(B) \) into two regions, with \( C_{\mathcal{P}(B)}^{1} \) on one side and \( C_{\mathcal{P}(B)}^{0} \) on the other side, is the union of these Voronoi facets. Each facet can be defined by an affine function over a compact subset of \( \mathbb{R}^{n-1} \), and the decision boundary is locally described by one of these functions.

Let \( \{ \epsilon_{j} \}_{j=1}^{n} \) be the canonical orthonormal basis of the vector space \( \mathbb{R}^n \). For \( y \in \mathbb{R}^n \), the \( j \)-th coordinate is \( y^j = y \cdot e_j \). Denote \( \tilde{y} = (y_2, ..., y_n) \in \mathbb{R}^{n-1} \) and let \( \mathcal{H} \{ h_j \} \) be the set of affine functions involved in the decision boundary. The affine boundary function \( h_j : \mathbb{R}^{n-1} \to \mathbb{R} \) is

\[
h_j(\tilde{y}) = y^1 = \left( p_j - \sum_{k \neq j} y_k v_j^k \right) / v_j^1, \quad (2)
\]

For the sake of simplicity, in the sequel \( h_j \) shall denote the function defined in (2) or its associated hyperplane \( \{ y \in \mathbb{R}^n : y \cdot v_j - p_j = 0 \} \) depending on the context.

Theorem 1. Consider a lattice defined by a VR basis \( B = \{ b_i \}_{i=1}^n \). Suppose that the \( n-1 \) points \( B \setminus \{ b_1 \} \) belong to the hyperplane \( \{ y \in \mathbb{R}^n : y \cdot e_1 = 0 \} \). Then, the decision boundary is given by a CPWL function \( f : \mathbb{R}^{n-1} \to \mathbb{R} \), expressed as

\[
f(\tilde{y}) = \bigwedge_{m=1}^M \left\{ \vee_{k=1}^n g_{m,k}(\tilde{y}) \right\}, \quad (3)
\]

where \( g_{m,k} \in \mathcal{H} \), \( 1 \leq l_m < \tau f \), and \( 1 \leq M \leq 2^n - 1 \).
Corollary 1. Consider a lattice defined by a VR basis \( B = \{ b_i \}_{i=1}^n \). Without loss of generality, assume that \( b_1 > 0 \). Suppose also that \( x^1 > \lambda^1, \forall x \in C_{P(B)}^1 \) and \( \forall \lambda \in T_f(x) \cap C_{P(B)}^0 \). Then, the decision boundary is given by a CPWL function as in (3).

See [15] - VII-B for the proofs. Some interesting lattices may not admit a VR basis, e.g. see \( E_6 \) in [5]. In this case, if \( \text{Vol}(P(B) \setminus \cup_{x \in C(B)} \nu(x)) \ll \text{Vol}(P(B)) \) then HLD yields efficient decoding. A basis satisfying this condition is called quasi-Voronoi-reduced. The new definition below presumes that \( B \) is quasi-Voronoi-reduced in order to make a successful discrimination of \( z_1 \) via the boundary function. Also, a surface in \( \mathbb{R}^n \) defined by a function \( g \) of \( n-1 \) arguments is written as \( \text{Surf}(g) = \{ (g(y), y) \in \mathbb{R}^n : y \in \mathbb{R}^{n-1} \} \).

Definition 2. Let \( B \) be a basis of \( \Lambda \). Assume that \( B \) and \( \{ e_i \}_{i=1}^n \) have the same orientation as in Theorem 1. The basis is called semi-Voronoi-reduced (SVR) if there exists at least two points \( x_1, x_2 \in C_{P(B)}^1 \) such that \( \text{Surf}(\bigvee_{k=1}^{\ell_1} g_{1,k}) \cap \text{Surf}(\bigvee_{k=1}^{\ell_2} g_{2,k}) \neq \emptyset \), where \( \ell_1, \ell_2 \geq 1, g_{1,k} \) are the facets between \( x_1 \) and all points in \( T_f(x_1) \cap C_{P(B)}^0 \), and \( g_{2,k} \) are the facets between \( x_2 \) and all points in \( T_f(x_2) \cap C_{P(B)}^0 \).

The above definition of a SVR basis imposes that the boundaries around two points of \( C_{P(B)}^1 \), defined by the two convex functions \( \bigvee_{k=1}^m g_{m,k}, m = 1, 2 \), have a non-empty intersection. Consequently, the min operator \( \wedge \) leads to a boundary function as in (3).

Corollary 1. \( P(B) \) for a SVR basis \( B \) admits a decision boundary defined by a CPWL function as in (3).

Example 1. Consider the lattice \( A_3 \) defined by the Gram matrix (4). To better illustrate the symmetries we rotate the basis to have \( b_1 \) collinear with \( e_1 \). Theorem 2 ensures that the decision boundary is a function. The function is illustrated on Figure 1 and its equation is (we omit the \( \tilde{y} \) in the formula to lighten the notations):

\[
f = [h_{p_1} \lor h_1 \lor h_2] \land [(h_{p_2} \lor h_1) \land (h_{p_2} \lor h_2)] \land [h_{p_3}],
\]

where \( h_{p_1}, h_{p_2} \) and \( h_{p_3} \) are hyperplanes orthogonal to \( b_1 \) (the \( p \) index stands for plateau). On Figure 2 each edge is orthogonal to a local affine function of \( f \) and labeled accordingly. The \( [\cdot] \) groups all the set of convex pieces of \( f \) that includes the same \( h_{p_j} \). Functions for higher dimensions are available in [15] - VII-A.

From now on, the default orientation of the basis with respect to the canonical axes of \( \mathbb{R}^n \) is assumed to be the one of Theorem 1. We call \( f \) the decision boundary function. The domain of \( f \) (its input space) is \( D \subset \mathbb{R}^{n-1} \). The domain \( D \) is the projection of \( P(B) \) on the hyperplane \( \{ e_i \}_{i=2}^n \). It is a bounded polyhedron that can be partitioned into disjunct convex (and thus connected) regions which we call linear regions. For any \( \tilde{y} \) in one of these regions, \( f \) is described by a unique local affine function \( h_j \). The number of those regions is equal to the number of affine pieces of \( f \).

V. FOLDING-BASED NEURAL DECODING OF \( A_n \)

In this section, we first prove that the lattice basis from the \( n \times n \) Gram matrix in (4) is VR (all bases are equivalent modulo rotations and reflections). We count the number of pieces of the decision boundary function. We then build a deep ReLU network which computes efficiently this function via folding. Finally, we use the fact that the \( n-2 \)-dimensional hyperplanes partitioning \( D \) are not in “general position” in order to prove that a two-layer network needs an exponential number of neurons to compute the function.

Consider a basis for the lattice \( A_n \) with all vectors from the first lattice shell. Also, the angle between any two basis vectors is \( \pi/3 \). Let \( J_n \) denote the \( n \times n \) all-ones matrix and \( I_n \) the identity matrix. The Gram matrix is

\[
\Gamma = GG^T = J_n + I_n.
\]

Theorem 3. A lattice basis defined by the Gram matrix (4) is Voronoi-reduced.
See [15] - VII-C for the proof. Consequently, discriminating a point in \( P(B) \) with respect to the decision boundary leads to an optimal decoder.

A. Number of pieces of the decision boundary function

We count the number of pieces, and thus linear regions, of the decision boundary function \( f \). We start with the following lemma involving \( i \)-simplices.

**Lemma 1.** Consider an \( A_n \)-lattice basis defined by the Gram matrix (4). The decision boundary function \( f \) has a number of affine pieces equal to

\[
\sum_{i=0}^{n} i \times (\# \text{ regular } i\text{-simplices}),
\]

(5)

where, for each \( i \)-simplex, only one corner \( x \) belongs to \( C_{P(B)}^0 \) and the other corners constitute the set \( \tau_f(x) \cap C_{P(B)}^0 \).

**Proof.** A key property of this basis is

\[
\forall x \in C_{P(B)}^0, \ x' \in A_n \setminus \{b_j, 0\}, \ 2 \leq j \leq n : x + b_j \in \tau_f(x + b_1), \ x + x' \not\in \tau_f(x + b_1) \cap C_{P(B)}^0.
\]

(6)

It is obvious that \( \forall x \in C_{P(B)}^0, x + b_1 \in \tau_f(x) \cap C_{P(B)}^0 \). This implies that any given point \( x \in C_{P(B)}^0 \) and its neighbors \( \tau_f(x) \cap C_{P(B)}^0 \) form a regular simplex \( S \) of dimension \( \frac{n(n+1)}{2} \). This clearly appears on Figure 2. Now, consider the decision boundary function of a \( k \)-simplex separating the top corner (i.e. \( C_{2}^k \)) from all the other corners (i.e. \( C_{n}^0 \)). This function is convex and has \( k \) pieces. The maximal dimensionality of such a simplex is obtained by taking the points \( 0, b_1, \) and the \( n-1 \) points \( b_j, j \geq 2 \).

**Theorem 4.** Consider an \( A_n \)-lattice basis defined by the Gram matrix (4). The decision boundary function \( f \) has a number of affine pieces equal to

\[
\sum_{i=1}^{n} i \cdot \binom{n-1}{n-i}.
\]

(7)

**Proof.** From Lemma 1, what remains to be done is to count the number of \( k \)-simplices. We walk in \( C_{P(B)}^0 \) and for each of the \( 2^{n-1} \) points \( x \in C_{P(B)}^0 \) we investigate the dimensionality of the simplex where the top corner is \( x + b_1 \). This is achieved by counting the number of elements in \( \tau_f(x + b_1) \cap C_{P(B)}^0 \), via the property given by (6). Starting from the origin, one can form a \( n \)-simplex with 0, \( b_1 \), and the \( n-1 \) other basis vectors. Then, from any \( b_j \), \( j \neq 1 \), one can only add the \( n-1 \) remaining basis vectors to generate a simplex in \( P(B) \). Indeed, if we add again \( b_j \), the point goes outside \( P(B) \). Hence, we get a \( n-1 \)-simplex and there are \( \binom{n-1}{2} \) ways to choose \( b_{j_1} \) any basis vectors except \( b_1 \). Similarly, if one starts the simplex from \( b_{j_1} + b_{j_2} \), one can form a \( n-2 \)-simplex in \( P(B) \) and there are \( \binom{n-2}{1} \) ways to choose \( b_{j_1} + b_{j_2} \). In general, there are \( \binom{n-1}{k} \) ways to form a \( n-k \)-simplex. Applying the previous lemma and summing over \( k = n-1 \) gives the announced result.

B. Decoding via folding

Obviously, at a location \( \tilde{y} \), we do not want to compute all affine pieces in (3) whose number is given by (7) in order to evaluate \( f \). To reduce the complexity of this evaluation, the idea is to exploit the symmetries of \( f \) by “folding” the function and mapping distinct regions of the input domain to the same location. If folding is applied sequentially, i.e. fold a region that has already been folded, it is easily seen that the gain becomes exponential. The notion of folding the input space in the context of neural networks was introduced in [9].

Given the basis orientation as in Theorem 1, the projection of \( b_j \) on \( D \) is \( b_j \) itself, for \( j \geq 2 \). We also denote the bisector hyperplane between two vectors \( b_j, b_k \) by \( BH(b_j, b_k) \) and its normal vector is taken to be \( v_{j,k} = b_j - b_k \). We define the folding transformation \( F : D \to D' \) as follows: let \( \tilde{y} \in D \), for all \( 2 \leq j < k \leq n \), compute \( \tilde{y} \cdot v_{j,k} \) (the first coordinate of \( v_{j,k} \) is zero). If the scalar product is non-positive, replace \( \tilde{y} \) by its mirror image with respect to \( BH(b_j, b_k) \). There exist \((n-1)(n-2)/2 \) hyperplanes for mirroring.

**Theorem 5.** Let us consider the lattice \( A_n \) defined by the Gram matrix (4). We have (i) \( D' \subset D \), (ii) for all \( \tilde{y} \in D \), \( f(\tilde{y}) = f(F(\tilde{y})) \) and (iii) \( f \) has exactly

\[
2n + 1
\]

pieces on \( D' \). This is to be compared with (7).

See [15] - VII-D for the proof.

**Example 1** (Continued). The function \( f \) restricted to \( D' \) (i.e. the function to evaluate after folding), say \( f_{D'} \), is

\[
f_{D'} = [h_{p1} \lor h_1] \land [h_{p2} \lor h_2] \land [h_{p3}].
\]

(9)

The general expression of \( f_{D'} \) for any dimension is available in [15] - VII-A.

C. From folding to a deep ReLU network

For sake of simplicity and without loss of generality, in addition to the standard ReLU activation function \( \text{ReLU}(a) = \max(0,a) \), we also allow the function \( \max(0,-a) \) and the identity as activation functions in the network.

To implement a reflection, one can use the following strategy. Step 1: rotate the axes to have the \( j \)-th axis \( e_j \) perpendicular to the reflection hyperplane and shift the point (i.e. the \( j \)-th coordinate) to have the reflection hyperplane at the origin. Step 2: take the absolute value of the \( j \)-th coordinate. Step 3: do the inverse operation of step 1.

Now consider the ReLU network illustrated in Figure 3. The edges between the input layer and the hidden layer represent the rotation matrix, where the \( j \)-th column is repeated twice, and \( p \) is a bias applied on the \( j \)-th coordinate. Within the dashed square, the absolute value of the \( j \)-th coordinate is computed and shifted by \( -p \). Finally, the edges between the hidden layer and the output layer represent the inverse rotation matrix. This ReLU network computes a reflection. We call it a reflection block.

All reflections can be naively implemented by a simple concatenation of reflection blocks. This leads to a very deep
and narrow network of depth $O(n^2)$ and width $O(n)$.

Regarding the $2n+1$ remaining pieces after folding, we have two options (in both cases, the number of operations involved is negligible compared to the previous folding operations). To
directly discriminate the point with respect to $f$, we implement
the HLD on these remaining pieces with two additional hidden
layers (see e.g. Figure 2 in [5]): project $y_{f,old}$ on the
$2n+1$ hyperplanes (with one layer of width $2n+1$) and
compute the associated Boolean equation with an additional
hidden layer. If needed, we can alternatively evaluate $f(y)$ via
$O(\log(n))$ additional hidden layers. First, compute the $n-1$ 2-
$\vee$ via two layers of size $O(n)$ containing several “max ReLU
networks” (see e.g. Figure 3 in [2]). Then, compute the $n-$
$\wedge$ via $O(\log(n))$ layers. Note that $f(y)$ can also be used for
discrimination via the sign of $y - f(y)$.

The total number of parameters in the whole network is
$O(n^4)$. In [15] - VII-G, we quickly discuss whether or not
this can be improved.

Eventually, the CPV is solved by using $n$ such networks in
parallel (this could also be optimized). The final network has
width $O(n^2)$ and depth $O(n^2)$.

D. Decoding via a shallow network

A two-layer ReLU network with $n$ inputs and $w_1$ neurons
in the hidden layer can compute a CPWL function with at
most $\sum_{j=0}^{w_1} \binom{n_j}{n}$ pieces [10]. This result is easily understood
by noticing that the non-differentiable part of $\max(a, 0)$ is a
$n-2$-dimensional hyperplane that separates two linear regions.
If one sums $w_1$ functions $\max(0, d_j y)$, where $d_j$, $1 \leq j \leq w_1$, is
a random vector, one gets $w_1$ of such $n-2$-hyperplanes.
The rest of the proof consists in counting the number of linear
regions that can be generated by these $w_1$ hyperplanes. The
number provided by the previous formula is attained if and
only if the hyperplanes are in general position. Clearly, in
our situation the $n-2$-hyperplanes partitioning $\mathcal{D}$ are not in
general position: the hyperplane arrangement is not simple.
The proof of the following theorem, available in [15] - VII-E,
consists in finding a lower bound on the number of such $n-2$
hyperplanes.

**Theorem 6.** A ReLU network with one hidden layer needs at
least
$$\sum_{i=2}^{n} (i - 1) \cdot \binom{n-1}{n-i}$$
neurons to solve the CVP for the lattice $A_n$. (10)

VI. CONCLUSIONS

We recently applied this theory to a SVR basis of lattices
$E_n$, $6 \leq n \leq 8$. The decision boundary function has a
number of pieces equal to
$$\sum_{i=3}^{n-1} \binom{n-3}{n-i} + 2 \left[ \frac{1}{2} + \frac{n-3}{2} \right]$$
$$+ \left[ \frac{1}{3} + \frac{n-3}{3} \right] + 1,$$
a number which we successfully linearized via folding.

From a learning perspective, our findings suggest that
many optimal decoders may be contained in the function
class $\Phi$ of deep ReLU networks. Learnability results of the
restricted model [1][3] show that the sample complexity is
then $m_{\Phi}(\epsilon, \delta) = O(WL \log(W)/\epsilon)$, avoiding the $1/\epsilon^2$ of
the general model (where $W$ is the number of parameters in
the network and $L$ the number of layers).

Additionally, the folding approach suits very well the non-
uniform finite sample bounds of the information bottleneck
framework [13]. Indeed, if we model the input of the network and
its $i$-th layer after the $i$-th reflection block by random
variables $Y$ and $Y_i$, clearly $I(Y;Y_i)$ is reduced compared to
$I(Y; Y_{i-1})$ for any distribution of $Y$.

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[15] The long version of this paper including an appendix is available at: www.josephboulos.com/thesis.pdf.