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Representation of hypergeometric products of higher nesting depths in difference rings

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Abstract

A non-trivial symbolic machinery is presented that can rephrase algorithmically a finite set of nested hypergeometric products in appropriately designed difference rings. As a consequence, one obtains an alternative representation in terms of one single product defined over a root of unity and nested hypergeometric products which are algebraically independent among each other. In particular, one can solve the zero-recognition problem: the input expression of nested hypergeometric products evaluates to zero if and only if the output expression is the zero expression. Combined with available symbolic summation algorithms in the setting of difference rings, one obtains a general machinery that can represent (and simplify) nested sums defined over nested products.

Keywords: difference rings, nested hypergeometric products, constant field, ring of sequences, zero recognition, algebraic independence, roots of unity products

1. Introduction

An important problem in symbolic summation is the simplification of sums defined over products to expressions in terms of simpler sums and products; in the best case, one might find an expression without sums. A first milestone was Gosper’s algorithm \cite{Gos78} and Zeilberger’s groundbreaking application for creative telescoping \cite{Zei91} where the summand is given by one hypergeometric product. Further extensions have been accomplished for $q$-hypergeometric and multibasic products \cite{PR97}, their mixed versions \cite{BP99} and ($q$-)multi-summation \cite{WZ92, Weg97, Ric03}. In addition, structural properties and further insight in this setting have been elaborated, e.g., in \cite{Pan95, CFFL11, CJKS13}. More generally, the holonomic system approach \cite{Zei90} and their refinements \cite{Chy00, Kon13, BRS18} represent, e.g., multi-sums over ($q$-)hypergeometric products by systems of linear recurrence relations.

In particular, Karr’s difference field approach \cite{Kar81, Kar85} paved the way for a general framework to represent rather complicated product expressions in a formal way. Here the generators of his $\Pi\Sigma$-field construction enables one to model (up to a certain level) indefinite nested products of the form

\begin{equation} P(n) = \prod_{k_1=\ell_1}^{n} f_1(k_1) \cdots \prod_{k_m=\ell_m}^{k_{m-1}} f_m(k_m) \tag{1} \end{equation}

where the multiplicands $f_i(n) = \frac{p_i(n)}{q_i(n)}$ for all $i$ with $1 \leq i \leq m$ are built by polynomial expressions $p_i(n)$ and $q_i(n)$ in terms of indefinite nested products that are again of the form (1). In Karr’s seminal works \cite{Kar81, Kar85}, which can be considered as the discrete version of Risch’s integration algorithm \cite{Ris69}, a sophisticated algorithm is provided that enables one to test if a given product

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\end{itemize}
representation (and sums defined over such products) are expressed properly in his \( \Pi \Sigma \)-field setting. As
a bonus, his toolbox and refinements in [Sch13, Sch07a, Sch08] enable one to decide constructively if a
given summand represented in a \( \Pi \Sigma \)-field has a solution in the same field or in an appropriate extension
of it. In addition, first contributions have been provided in [Sch05, AP10] to simplify products further
such that the degrees in the numerators and denominators are minimal.

For such complicated classes of products the following task is non-trivial: given an arbitrary ex-
pression in terms of nested products \([\square]\), design algorithmically an appropriate difference field or ring
in which the expression can be represented and in which one can solve, e.g., the (creative) telescop-
ning problem. Already for a hypergeometric product \( P_1(n) = \prod_{k=\ell}^{n} f(k) \) with a rational function
\( f(x) \in \mathbb{K}(x)^* \) and \( \ell \in \mathbb{Z}_{\geq 0} \) chosen properly (i.e., \( P_1(n) \) is well defined and nonzero for all \( n \in \mathbb{Z}_{\geq 0} \)), it
has been shown in [Sch05] that Karr’s \( \Pi \Sigma \)-fields are not sufficient: namely, such a product \( P_1(n) \) can
be represented in a \( \Pi \Sigma \)-field if and only if it cannot be rewritten in the form \( P_1(n) = \zeta^n r(n) \) where
\( r(x) \in \mathbb{K}(x) \) and \( \zeta \in \mathbb{K} \) is a primitive \( \lambda \)-th root of unity with \( \lambda > 1 \). More precisely, such objects

Motivated by this observation (among others) a refined difference ring theory has been elaborated
in [Sch16, Sch17] that combines big parts of Karr’s general framework together with generators of the form \( \zeta^n \). In this way, not only the hypergeometric product \( P_1(n) \), but more generally any polynomial
expression in terms of hypergeometric products can be represented in the class of so-called \( \Pi \Sigma \)-
extensions. The first algorithms derived in [Sch05, Sch14] require that the input products are defined
over \( \mathbb{K}(x) \) where \( \mathbb{K} = \mathbb{Q}(\kappa_1, \ldots, \kappa_u) \) with \( u \geq 0 \) is a rational function field defined over the rational
numbers \( \mathbb{Q} \). More generally, a complete algorithm has been elaborated in [OS18] (utilizing ideas
from [Ge93, Sch14]) that can represent a finite set of hypergeometric products over \( \mathbb{K}(x) \) where \( \mathbb{K} =
K(y_1, \ldots, y_r) \) is a rational function field defined over an algebraic number field \( K \). In addition, using
algorithms from [BP99] it can deal also with \( q \)-hypergeometric, multibasic and mixed hypergeometric
products. Finally, a general framework has been elaborated in [Sch20] that considers single nested
products defined over a general class of difference fields; an extra bonus is that the latter approach
constructs a difference ring in which the given products are rephrased optimally with following property:
the number of generators of the ring and the order \( \lambda \) of the used \( \zeta^n \) are minimal.

A remarkable feature of the above algorithms [Sch05, Sch14, OS18, Sch20] is that the given input
expression of hypergeometric products (and their generalized versions) are rephrased in terms of a finite
set of alternative products \( Q_1(n), \ldots, Q_s(n) \) together with a distinguished root of unity product \( \zeta^n \)
such that the sequences produced by \( Q_1(n), \ldots, Q_s(n) \) are algebraically independent among each other.
For this result we rely on ideas of [Sch17] that are inspired by [Sch05, HS08]; compare also [CFLL11].
We remark further that these results are also connected to [KZ08] that can compute all algebraic
relations of \( C \)-finite solutions (i.e., solutions of homogeneous recurrences with constant coefficients).
We emphasize that the algorithms from [Sch05, AP10, Sch14, OS18, Sch20] can be utilized to simplify
hypergeometric solutions [Pet92, Hoe99, ABPS20] of linear difference equations and can be combined
with symbolic summation algorithms [Sch04a, Sch15, Sch08] to simplify more general solutions, such as
d’Alambertian solutions [AP94, AZ96] and Liouvillian solutions [HS99, PZ13].

In this article we aim at extending this toolbox significantly for the general class of nested hyper-
geometric products that can be defined as follows.

**Definition 1.** Let \( \mathbb{K}(x) \) be a rational function field\(^1\) and let \( f_1(x), \ldots, f_m(x) \in \mathbb{K}(x)^* \). Furthermore,
let \( \ell_1, \ldots, \ell_m \in \mathbb{Z}_{\geq 0} \) such that for all \( i \) with \( 1 \leq i \leq m \), \( f_i(j) \) is non-zero and has no pole for all \( j \in \mathbb{Z}_{\geq 0} \)
with \( j \geq \ell_i \). Then the indefinite product expression \([\square]\) is called a hypergeometric product in \( n \) of nesting
depth \( m \). The vector \( (f_1(x), \ldots, f_m(x)) \in (\mathbb{K}(x)^*)^m \) is also called the multiplicand representation
of \( P(n) \). If \( f_i(x) \in \mathbb{K}^* \) for \( 1 \leq i \leq m \), then we call \([\square]\) a constant or geometric product in \( n \) of nesting
depth \( m \). Further, we define the set of ground expressions with\(^2\) \( \mathbb{K}(n) = \{ f(n) \mid f(x) \in \mathbb{K}(x) \} \).

\(^1\)Throughout this article all fields and rings have characteristic 0.

\(^2\)Their elements are considered as expressions that can be evaluated for sufficiently large \( n \in \mathbb{Z}_{\geq 0} \).
define $\text{Prod}_n(G)$ with $G \subseteq \mathbb{K}(x)$ as the set of all such products where the multiplicant representations are taken from $G$. Furthermore, we introduce the set of product monomials $\text{Prod}_n(G)$ as the set of all elements

$$a(n)P_1(n)^{\nu_1} \cdots P_e(n)^{\nu_e}$$

with $a(x) \in G$, $e \in \mathbb{Z}_{\geq 0}$, $\nu_1, \ldots, \nu_e \in \mathbb{Z}$ and $P_1(n), \ldots, P_e(n) \in \text{Prod}_n(G)$. Finally, we introduce the set of product expressions $\text{Prod}E_n(G)$ as the set of all elements

$$A(n) = \sum_{\nu=(\nu_1, \ldots, \nu_e) \in S} a_{\nu}(n) P_1(n)^{\nu_1} \cdots P_e(n)^{\nu_e}$$

(2)

with $e \in \mathbb{Z}_{\geq 0}$, $S \subseteq \mathbb{Z}^e$ finite, $a_\nu(x) \in G$ for $\nu \in S$ and $P_1(n), \ldots, P_e(n) \in \text{Prod}_n(G)$. Note that $\text{Prod}_n(G) \subseteq \text{Prod}_n(G) \subseteq \text{Prod}E_n(G)$.

Utilizing the available algorithms from \cite{OS18} we will obtain enhanced algorithms that can rephrase expressions from $\text{Prod}E_n(G(x))$ in the setting of $\text{RII}$-extensions. As a consequence we will solve the following problem; for further details see Theorem 9 and Corollary 4 below.

**Problem RPE: Representation of Product Expressions.**

Let $\mathbb{K} = K(\kappa_1, \ldots, \kappa_u)$ be a rational function field with $e \geq 0$ over an algebraic number field $K$. Given $A(n) \in \text{Prod}E_n(\mathbb{K}(x))$. Find $B(n) \in \text{Prod}E_n(\mathbb{K}(x))$ with $\mathbb{K} = \tilde{K}(\kappa_1, \ldots, \kappa_u)$ where $\tilde{K}$ is an algebraic field extension of $K$, and a non-negative integer $\delta \in \mathbb{Z}_{\geq 0}$ with the following properties:

1. $A(n) = B(n)$ for all $n \in \mathbb{Z}_{\geq 0}$ with $n \geq \delta$;
2. All the products $P_1(n), \ldots, P_e(n) \in \text{Prod}_n(\mathbb{K}(x))$ arising in $B(n)$ (apart from the distinguished product $\zeta^n$ with $\zeta$ a root of unity) are algebraically independent among each other.
3. The zero-recognition property holds, i.e., $A(n) = 0$ holds for all $n$ from a certain point on if and only if $B(n)$ is the zero-expression.

The full machinery have been implemented within Ocansey’s Mathematica package $\text{NestedProducts}$ whose functionality will be illustrated in Section 6.3 below; for additional aspects we refer also to \cite{Oca19}. We expect that this implementation will open up new applications, e.g., in combinatorics, such as non-trivial evaluations of determinants \cite{MRRJS93, Zei96, Kra01}. In particular, in interaction with the symbolic summation algorithms available in the package $\text{Sigma}$ \cite{Sch07} one obtains a fully automatic toolbox to tackle nested sums defined over nested hypergeometric products.

The outline of the article is as follows. In Section 2 we will introduce rewrite rules that enable one to transform expressions from $\text{Prod}E_n(\mathbb{K}(x))$ to a more suitable form (see Proposition 2 below) to solve Problem RPE. Given this tailored form, we show in Section 3 how such expressions can be rephrased straightforwardly in terms of multiple-chain AP-extensions. In order to solve Problem RPE we have to refine this difference ring construction. Namely, in Section 4 we introduce RII-extensions: these are AP-extensions where during the construction the set of constants remain unchanged. In particular, we will elaborate that such rings can be straightforwardly embedded into the ring of sequences and will provide structural theorems that will prepare the ground to solve Problem RPE. With these results we will present in Section 5 the main steps how nested products can be represented in RII-extensions. In Section 6 we will combine all these ideas yielding a complete algorithm for Problem RPE that is summarized in Theorem 9 and Corollary 4. In addition, we will illustrate with non-trivial examples how one can solve Problem RPE with the new Mathematica package $\text{NestedProducts}$. The conclusions are given in Section 7.
2. Preprocessing hypergeometric products of finite nesting depth

In order to support our machinery to solve \textbf{Problem RPE}, the arising products $P(n)$ in $A(n) \in \text{ProdE}_n(\mathbb{K}(x))$ (e.g., given in (2)) will be transformed to a particularly nice form. We will illustrate each preprocessing step with an example and then summarize the derived result in Proposition 2 below.

Let $\mathbb{K}(x)$ be a rational function field together with the zero-function (in short $Z$-function) defined by

$$Z(p) = \max\{(k \in \mathbb{Z}_{\geq 0} \mid p(k) = 0)\} + 1 \text{ for any } p \in \mathbb{K}[x]$$

with $\max(\emptyset) = -1$. We call $\mathbb{K}$ computable if all basic field operations are computable. Note that if $\mathbb{K}$ is a rational function field over an algebraic number field, then $\mathbb{K}$ and also its $Z$-function are computable.

We start with the hypergeometric product in $n$ of nesting depth $m \in \mathbb{Z}_{\geq 0}$ given by

$$P(n) = \prod_{k_1=\ell_1}^{\ell_1} f_1(k_1) \prod_{k_2=\ell_2}^{\ell_2} f_2(k_2) \cdots \prod_{k_m=\ell_m}^{\ell_m} f_m(k_m) \in \text{Prod}_n(\mathbb{K}(x))$$

where $f_i(x) \in \mathbb{K}(x)^*$ and $\ell_i \in \mathbb{Z}_{\geq 0}$ for all $1 \leq i \leq m$. Note that by definition $P(n) \neq 0$ for all $n \in \mathbb{Z}_{\geq 0}$.

In particular, no poles arise for any evaluation at $n \in \mathbb{Z}_{\geq 0}$. We remark that the $Z$-function can be used to specify the lower bounds $\ell_i \in \mathbb{Z}_{\geq 0}$ such that this property holds. Then $P(n)$ in $\text{Prod}_n(\mathbb{K}(x))$ is preprocessed as follows.

2.1. Transformation of indefinite products to product factored form

The first transformation is based on the following simple observation.

\textbf{Proposition 1.} For $P(n)$ given in (4) with multiplicands $f_1, \ldots, f_m \in \mathbb{K}(x)^*$ we have

$$P(n) = \left( \prod_{k_1=\ell_1}^{\ell_1} f_1(k_1) \right) \left( \prod_{k_2=\ell_2}^{\ell_2} f_2(k_2) \right) \cdots \left( \prod_{k_m=\ell_m}^{\ell_m} f_m(k_m) \right) \in \text{Prod}_n(\mathbb{K}(x)).$$

\textbf{Definition 2.} The right hand side of (5) is also called a product factored form of $P(n)$. Moreover, a product of the form

$$P'(n) = \prod_{k_1=\ell_1}^{\ell_1} \cdots \prod_{k_m=\ell_m}^{\ell_m} p_m(k_m)$$

is also called a product in factored form. In particular, we also call $p_m(x) \in \mathbb{K}(x)^*$ (instead of $(1, \ldots, 1, p_m(x))$) the multiplicand representation of $P'(n)$.

Further, for $1 \leq i \leq m$ write

$$f_i = u_i f_{i,1}^{e_{i,1}} \cdots f_{i,r_i}^{e_{i,r_i}} \in \mathbb{K}(x)$$

in its complete factorization. This means that $f_i$ can be decomposed by $u_i \in \mathbb{K}^*$ and irreducible monic polynomials $f_{i,j} \in \mathbb{K}[x] \setminus \mathbb{K}$ with $e_{i,j} \in \mathbb{Z}$ for some $1 \leq j \leq r_i$ with $r_i \in \mathbb{Z}_{\geq 0}$. Substituting (6) into the right-hand side of (5) and expanding the product quantifiers over each factor in (6) we get

$$P(n) = A_1(n) A_2(n) \cdots A_m(n) \in \text{Prod}_n(\mathbb{K}(x))$$

where

$$A_i(n) = \left( \prod_{k_1=\ell_1}^{\ell_1} \cdots \prod_{k_i=\ell_i}^{\ell_i} u_i \right) \left( \prod_{k_1=\ell_1}^{\ell_1} \cdots \prod_{k_i=\ell_i}^{\ell_i} f_{i,1}(k_i) \right)^{e_{i,1}} \cdots \left( \prod_{k_1=\ell_1}^{\ell_1} \cdots \prod_{k_i=\ell_i}^{\ell_i} f_{i,r_i}(k_i) \right)^{e_{i,r_i}}$$

for all $1 \leq i \leq m$. In particular, the first product on the right hand side in (7) with innermost multiplicand $u_i \in \mathbb{K}^*$ is a geometric product of nesting depth $i$ in $\text{Prod}_n(\mathbb{K})$, while the rest are nesting depth $i$ hypergeometric products in $\text{Prod}_n(\mathbb{K}(x))$ which are not geometric products.
Example 1. Let $\mathbb{K} = \mathbb{Q}(\sqrt{3})$ and $\mathbb{K}(x)$ be the rational function field over $\mathbb{K}$ with the $Z$-function $\mathbb{Z}$.

Suppose we are given the nesting depth 2 hypergeometric product

$$P(n) = \prod_{k=1}^{n} \frac{24k + 1}{-\sqrt{3}} \prod_{j=3}^{k} \frac{-2(j^3 - 3j + 2)}{5(j^2 - j - 2)} \in \text{Prod}_n(\mathbb{K}(x)).$$

(8)

Then with

$$A_1(n) = \left( \prod_{k=1}^{n} \frac{1}{\sqrt{3}} \right) \left( \prod_{k=1}^{n} \delta \right) \left( \prod_{k=1}^{n} k \right)$$

$$A_2(n) = \left( \prod_{k=1}^{n} \frac{k}{\sqrt{3}} \right) \left( \prod_{k=1}^{n} k \right) \left( \prod_{k=1}^{n} k \right)$$

(9)

(10)

equation (8) can be written in the form

$$P(n) = A_1(n) A_2(n) \in \text{Prod}_n(\mathbb{K}(x))$$

where the multiplicand representations of the products in $A_1(n)$ and $A_2(n)$ are either from $\mathbb{K}$ or are irreducible polynomials from $\mathbb{K}[x]$.

2.2. Synchronization of lower bounds

Another transformation will guarantee that all arising products have the same lower bound, i.e., that the expression is $\delta$-refined for some $\delta \in \mathbb{Z}_{\geq 0}$.

Definition 3. Let $\mathbb{K}(x)$ be a rational function field over a field $\mathbb{K}$ and $\delta \in \mathbb{Z}_{\geq 0}$. $H(n) \in \text{Prod}_n(\mathbb{K}(x))$ is said to be $\delta$-refined if the lower bounds in all the arising products of $H(n)$ are $\delta$.

Such a transformation of a given product expression to a $\delta$-refined version can be accomplished by taking $\delta$ to be the maximum of all arising lower bounds within the given expression.

Example 2 (Cont. Example 1). In $P(n)$ (resp. $A_1(n)$ and $A_2(n)$ of Example 1) we choose $\delta = 3$.

Namely, for all $1 \leq i \leq m$, rewrite each product in (7) such that the lower bounds are synchronized to $\delta$. More precisely we apply the formula

$$\prod_{k_1=\ell_1}^{n} \prod_{k_2=\ell_2}^{k_{i-1}} h(k_i) = \left( \prod_{k_1=\ell_1}^{n} \prod_{k_2=\ell_2}^{\delta} h(k_i) \right) \left( \prod_{k_1=\ell_1}^{n} \prod_{k_2=\ell_2}^{\delta} h(k_i) \right)$$

(11)

to each of the products in (7). Note that the first product on the right-hand side in (11) evaluates to a constant in $\mathbb{K}^*$, the last product is from $\text{Prod}_n(\mathbb{K}(x))$, and all the remaining products (after all finite multiplications are carried out) are from $\text{Prod}_n(\mathbb{K})$. Summarizing we obtain

$$\hat{P}(n) = \hat{A}_1(n) \hat{A}_2(n) \cdots \hat{A}_m(n) \in \text{Prod}_n(\mathbb{K}(x))$$

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with 
\[
\tilde{A}_i(n) = a_i \left( \prod_{k_1=\delta}^n \bar{u}_{i,1} \right) \cdots \left( \prod_{k_1=\delta}^n \bar{u}_{i,i} \right) \left( \prod_{k_1=\delta}^n f_{i,r_i}(k_i) \right)^{e_{i,r_i}}
\]  
(12)

where \( a_i, \bar{u}_{i,j} \in \mathbb{K}^* \) for some \( j \in \mathbb{Z}_{\geq 0} \). Since \( \delta \) is chosen as the maximum among all lower bounds of the input expression, no poles or zero-evaluations will be introduced. As a consequence, the obtained result is again an element from \( \text{ProdM}_n(\mathbb{K}(x)) \). In particular, we have that \( A_i(n) = \tilde{A}_i(n) \) for all \( n \geq \max(0, \delta - 1) \) and consequently, \( P(n) = \tilde{P}(n) \) holds for all \( n \geq \max(\delta - 1, 0) \).

Example 3 (Cont. Example 2). Synchronizing the lower bounds of each product factor in (9) and (10) to 3 computed in Example 2 and rewriting each product factor in (9) and (10) we get

\[
\tilde{A}_1(n) = \frac{1225}{3} \left( \prod_{k=3}^n -1 \right) \left( \prod_{k=3}^n \sqrt{3} \right)^{-1} \left( \prod_{k=3}^n 24 \right) \left( \prod_{k=3}^n (k + \frac{1}{2}) \right),
\]  
(13)

\[
\tilde{A}_2(n) = \left( \prod_{k=3}^n k \right)^{-1} \left( \prod_{k=3}^n 5 \right)^{-1} \left( \prod_{k=3}^n 2 \right)^{2} \left( \prod_{k=3}^n (j - 2) \right)^{2} \left( \prod_{k=3}^n (j - 1) \right)^{-1} \left( \prod_{k=3}^n (j + 1) \right)^{-1} \left( \prod_{k=3}^n (j + 2) \right). 
\]  
(14)

In particular, for \( i = 1, 2 \), and for all \( n \geq \delta - 1 \) where \( \delta = 3 \), \( A_i(n) = \tilde{A}_i(n) \) holds. Consequently, with \( P(n) = \tilde{A}_1(n) \tilde{A}_2(n) \) we have that \( P(n) = \tilde{P}(n) \) holds for all \( n \geq 2 \).

Since geometric products never introduce poles or zeroes, we can bring each geometric product in (12) to a 1-refined form by using a similar formula as given in (11). This yields

\[
P'(n) = A'_1(n) A'_2(n) \cdots A'_n(n) \in \text{ProdM}_n(\mathbb{K}(x))
\]

where

\[
A'_1(n) = \tilde{a}_i \left( \prod_{k_1=1}^n \tilde{u}_{i,1} \right) \cdots \left( \prod_{k_1=1}^n \tilde{u}_{i,i} \right) \left( \prod_{k_1=\delta}^n f_{i,r_i}(k_i) \right)^{e_{i,r_i}}
\]  
(15)

with \( \tilde{a}_i, \tilde{u}_{i,j} \in \mathbb{K}^* \) for some \( j \in \mathbb{Z}_{\geq 0} \). In particular we have that, \( A_i(n) = A'_i(n) \) holds for all \( n \geq \max(\delta - 1, 0) \) and consequently, \( P(n) = P'(n) \) holds for all \( n \geq \max(0, \delta - 1) \). By rearranging the arising products in \( P'(n) \) we obtain the decomposition

\[
P'(n) = c G(n) H(n)
\]  
(16)

with \( c \in \mathbb{K}^*, \ G(n) \in \text{ProdM}_n(\mathbb{K}) \) is composed multiplicatively by geometric products in factored form of nesting depth at most \( m \) which are 1-refined and \( H(n) \in \text{ProdM}_n(\mathbb{K}(x)) \) is composed multiplicatively by hypergeometric products (which are not geometric) in factored form of nesting depth at most \( m \) which are \( \delta \)-refined. In particular, the multiplicant representations are given by monic irreducible polynomials.
Example 4 (Cont. Example 3). Synchronizing the lower bounds of each geometric product in (13) and (14) to 1 and rewriting these geometric products we get

\[ A'_1(n) = \frac{1225}{576} \left( \prod_{k=1}^{n} \frac{1}{k} \right) \left( \prod_{k=1}^{n} \frac{\sqrt{3}}{k} \right)^{-1} \left( \prod_{k=1}^{n} \frac{24}{k} \right) \left( \prod_{k=3}^{n} \frac{k + \frac{1}{24}}{k} \right), \]

\[ A'_2(n) = -\frac{2}{5} \left( \prod_{k=1}^{n} \frac{4}{k} \right)^{-1} \left( \prod_{k=1}^{n} \frac{25}{k} \right) \left( \prod_{k=1}^{n} \frac{k}{k=1}=1 \right) \left( \prod_{k=1}^{n} \frac{5}{k=1}=1 \right)^{-1} \left( \prod_{k=3}^{n} \frac{(j - 2)}{k=3}=3 \right) \left( \prod_{k=3}^{n} \frac{(j + 1)}{k=3}=3 \right)^{-1} \left( \prod_{k=3}^{n} \frac{(j + 2)}{k=3}=3 \right). \]

In particular, for \( i = 1, 2 \), and for all \( n \geq 2 \), \( A_i(n) = A'_i(n) \) holds. In total we obtain

\[ P'(n) = A'_1(n) A'_2(n) = c G(n) H(n) \]

with

\[ c = -\frac{245}{288}, \]

\[ G(n) = \left( \prod_{k=1}^{n} \frac{1}{k} \right) \left( \prod_{k=1}^{n} \frac{\sqrt{3}}{k} \right)^{-1} \left( \prod_{k=1}^{n} \frac{24}{k} \right) \left( \prod_{k=1}^{n} \frac{25}{k} \right) \left( \prod_{k=1}^{n} \frac{k}{k=1}=1 \right) \left( \prod_{k=1}^{n} \frac{5}{k=1}=1 \right)^{-1} \left( \prod_{k=3}^{n} \frac{k}{k=3}=3 \right) \left( \prod_{k=3}^{n} \frac{k}{k=3}=3 \right)^{-1} \left( \prod_{k=3}^{n} \frac{k}{k=3}=3 \right), \]

\[ H(n) = \left( \prod_{k=3}^{n} \frac{(k + \frac{1}{24})}{k=3} \right) \left( \prod_{k=3}^{n} \frac{k}{k=3} \right)^{-1} \left( \prod_{k=3}^{n} \frac{(j - 2)}{k=3} \right) \left( \prod_{k=3}^{n} \frac{(j + 1)}{k=3} \right)^{-1} \left( \prod_{k=3}^{n} \frac{(j + 2)}{k=3} \right). \]

such that \( P(n) = P'(n) \) holds for all \( n \geq 2 \).

2.3. Shift-coprime representation

Finally, we turn our focus to the class of hypergeometric products given in factored form, and whose innermost multiplicands are irreducible monic polynomials. In order to reduce this class of products further, we will need the following definition.

Definition 4. Two nonzero polynomials \( f(x) \) and \( h(x) \) in the polynomial ring \( \mathbb{K}[x] \) are said to be shift-coprime if for all \( k \in \mathbb{Z} \) we have that \( \gcd(f(x), h(x + k)) = 1 \). Furthermore, \( f(x) \) and \( h(x) \) are called shift-equivalent if there is a \( k \in \mathbb{Z} \) such that \( \frac{f(x+k)}{h(x)} \in \mathbb{K} \).

It is immediate that the shift-equivalence in Definition 4 induces an equivalence relation on the set of all irreducible polynomials. Let \( D = \{ f_1, \ldots, f_e \} \subseteq \mathbb{K}[x] \) where all elements are irreducible and shift equivalent among each other. Then we call \( f_i \in D \) with \( i \in \{ 1, 2, \ldots, e \} \) the leftmost polynomial in \( D \) if for all \( h \in D \) there is a \( k \in \mathbb{Z}_{\geq 0} \) with \( \frac{f_i(x+k)}{h(x)} \in \mathbb{K} \). It is well known that \( \frac{f_i(x+k)}{h(x)} \in \mathbb{K} \) iff \( k \in \mathbb{Z} \) is a root of \( p(z) = \text{res}_x(f(x), f_i(x + z)) \in \mathbb{K}[z] \); compare [PWZ96] Sec. 5.3. In particular, if \( \mathbb{K} \) is computable and one can factorize univariate polynomials over \( \mathbb{K} \), one can determine all integer roots of \( p(z) \) and thus can decide constructively if there is a \( k \in \mathbb{Z} \) with \( \frac{f_i(x+k)}{h(x)} \in \mathbb{K} \). All the above properties (and slight generalizations) play a crucial role in symbolic summation; compare [Abr71, Pau95, Sch05, AP08, CFFL11]. In particular, the following simple lemma is heavily used within symbolic summation; see also [Sch05] Lemma 4.12.
Lemma 1. Let $\mathbb{K}(x)$ be a rational function field and let $f(x)$, $h(x) \in \mathbb{K}[x] \setminus \mathbb{K}$ be monic irreducible polynomials that are shift equivalent. Then there is a $g \in \mathbb{K}(x)^*$ with $h(x) = \frac{g(x+1)}{g(x)} f(x)$ were all the monic irreducible factors in $g$ are shift equivalent to $f(x)$ (resp. $h(x)$). If $\mathbb{K}$ is computable and one can factorize polynomials over $\mathbb{K}$, then such a $g$ can be computed.

Proof. Since $f(x)$ and $h(x)$ are shift equivalent and monic, there is a $k \in \mathbb{Z}$ with $f(x+k) = h(x)$. If $k \geq 0$, set $g := \prod_{i=0}^{k-1} f(x+i)$. Then

$$\frac{g(x+1)}{g(x)} = \frac{f(x+k)}{f(x)} = \frac{h(x)}{f(x)}.$$  

On the other hand, if $k < 0$, set $g := \prod_{i=1}^{-k} \frac{1}{f(x+i)}$. Then

$$\frac{g(x+1)}{g(x)} = \frac{1/f(x)}{1/f(x+k)} = \frac{f(x+k)}{f(x)} = \frac{h(x)}{f(x)}.$$  

By construction all irreducible monic factors in $g(x)$ are shift equivalent to $f(x)$. Furthermore, $k$ can be computed if $\mathbb{K}$ is computable and one can factorize polynomials over $\mathbb{K}$. \hfill \Box

Example 5. Let $\mathbb{K}(x)$ be a rational function field as defined in Example 1. Let $\mathcal{D}$ be the set defined by the multiplicand representations of the products in factored form given in [19]. That is, $\mathcal{D} = \{ f_1(x), f_2(x), \ldots, f_5(x) \} \subseteq \mathbb{K}[x] \setminus \mathbb{K}$ where

$$f_1(x) = x - 2, \quad f_2(x) = x - 1, \quad f_3(x) = x + 1, \quad f_4(x) = x + 2 \quad \text{and} \quad f_5(x) = x + \frac{1}{27}. \quad (20)$$

Since $f_1(x)$ is shift equivalent with $f_2(x), f_3(x), f_4(x)$, i.e.,

$$f_1(x+1) = f_2(x), \quad f_1(x+3) = f_3(x), \quad f_1(x+4) = f_4(x),$$

they fall into the same equivalence class $\mathcal{E}_1 = \{ f_1(x), f_2(x), f_3(x), f_4(x) \}$. The other equivalence class is $\mathcal{E}_2 = \{ f_5(x) \}$. For each of these equivalence classes $\mathcal{E}_1$ and $\mathcal{E}_2$, take their leftmost elements: $f_1(x)$ and $f_5(x)$ respectively. Then by Lemma 1 we can express the elements of each equivalence class in terms of the leftmost polynomial $f_1(x)$ or $f_5(x)$. More precisely, we have the following relations for the equivalence class $\mathcal{E}_1$:

$$f_2(x) = \frac{g_1(x+1)}{g_1(x)} f_1(x), \quad \text{with} \quad g_1(x) = (x-2), \quad (21)$$

$$f_3(x) = \frac{g_2(x+1)}{g_2(x)} f_1(x), \quad \text{with} \quad g_2(x) = (x-2)(x-1)x, \quad (22)$$

$$f_4(x) = \frac{g_3(x+1)}{g_3(x)} f_1(x), \quad \text{with} \quad g_3(x) = (x-2)(x-1)x(x+1). \quad (23)$$

Finally, we reduce each component of the hypergeometric product expression $H(n)$ given by [19]. We will begin with the nesting depth 2 hypergeometric products in factored form whose innermost multiplicand corresponds to the polynomial $f_4(x)$. Using [23] the product in factored form reduces as follows:

$$\prod_{k=3}^{n} f_a(j) = \prod_{k=3}^{n} \prod_{j=3}^{k} (j+2) = \left( \prod_{k=3}^{n} \prod_{j=3}^{k} g_a(j+1) \right) \left( \prod_{k=3}^{n} \prod_{j=3}^{k} f_a(j) \right) = \left( \prod_{k=3}^{n} \frac{1}{24} \right) \left( \prod_{k=3}^{n} (k-1) \right) \left( \prod_{k=3}^{n} (k+1) \right) \left( \prod_{k=3}^{n} (k+2) \right) \left( \prod_{k=3}^{n} (j-2) \right)$$

$$= 576 \left( \prod_{k=3}^{n} \frac{1}{24} \right) \left( \prod_{k=3}^{n} (k-1) \right) \left( \prod_{k=3}^{n} (k+1) \right) \left( \prod_{k=3}^{n} (k+2) \right) \left( \prod_{k=3}^{n} (j-2) \right). \quad (24)$$

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Using (22) and (21), a similar reduction can be achieved for the nesting depth 2 hypergeometric products in factored form arising in $H(n)$ whose innermost multiplicands correspond to the polynomials $f_3(x)$ and $f_2(x)$ respectively. In particular, we have the following:

\[
\prod_{k=3j=3}^{n} k f_3(j) = \prod_{k=3j=3}^{n} \prod_{k=3j=3}^{n} (j+1) = 36 \left( \prod_{k=1}^{n} \right) \left( \prod_{k=3}^{n} (k-1) \right) \prod_{k=3}^{n} k \left( \prod_{k=3}^{n} (k+1) \right) \prod_{k=3}^{n} \prod_{k=3j=3}^{n} (j-2), \tag{25}
\]

\[
\prod_{k=3j=3}^{n} k f_2(j) = \prod_{k=3j=3}^{n} \prod_{k=3j=3}^{n} (j-1) = \left( \prod_{k=3}^{n} (k-1) \right) \left( \prod_{k=3}^{n} \prod_{k=3j=3}^{n} (j-2) \right). \tag{26}
\]

**Remark 1.** Suppose we are given an expression $A(n) \in \text{ProdE}_n(\mathbb{K}(x))$ (e.g., given in [2]) in terms of $\delta$-refined hypergeometric products of finite nesting depth in factored form where all multiplicand representations are irreducible monic polynomials. Choose

\[
P'(n) = \prod_{k_1=\delta}^{n} \cdots \prod_{k_{m-1}=\delta}^{k_{m-2}} \prod_{k_m=\delta}^{k_{m-1}} h(k_m) \tag{27}
\]

from $A(n)$ with the multiplicand representation $h(x) \in \mathbb{K}[x]$. Furthermore, among all shift-equivalent multiplicand representations within the given product expression $A(n)$, let $f(x)$ be the leftmost polynomial which lies in the same equivalence class with $h(x)$. By assumption $f(x)$ and $h(x)$ are monic irreducible with $f(n) \neq 0$ and $h(n) \neq 0$ for all $n \geq \delta$. Take $k \in \mathbb{Z}_{\geq 0}$ with $h(x+k) = f(x)$. Then by Lemma 1 we can take $g := \prod_{i=0}^{k-1} f(x+i) \in \mathbb{K}[x]$ such that $h(x) = \frac{g(x+1)}{g(x)} f(x)$ holds. Thus

\[
P'(n) = \prod_{k_1=\delta}^{n} \cdots \prod_{k_{m-1}=\delta}^{k_{m-2}} \prod_{k_m=\delta}^{k_{m-1}} g(k_m + 1) f(k_m) = \prod_{k_1=\delta}^{n} \prod_{k_2=\delta}^{k_{m-1}} \prod_{k_m=\delta}^{k_{m-2}} g(k_m + 1) \prod_{k_m=\delta}^{k_{m-1}} g(\delta) f(k_m)
\]

\[
= \left( \prod_{k_1=\delta}^{n} \cdots \prod_{k_{m-1}=\delta}^{k_{m-2}} g(\delta) \right)^{-1} \left( \prod_{k_1=\delta}^{n} \cdots \prod_{k_{m-1}=\delta}^{k_{m-2}} g(k_m + 1) \right) \left( \prod_{k_1=\delta}^{n} \cdots \prod_{k_{m-1}=\delta}^{k_{m-2}} f(k_m) \right)
\]

\[
= G'(n) H'(n)
\]

Note that this reduction of a product of nesting depth $m$ leads to a new hypergeometric product $H'(n)$ in factored form of nesting depth less than $m$ where for the multiplicand representation $h'(x) := g(x+1)$ we have that $h'(n) \neq 0$ for all $n \geq \max(\delta - 1, 0)$. In particular, $h'(x)$ consists of monic irreducible factors which are again shift-equivalent to $f(x)$. In addition taking all these new factors together with $f(x)$, it follows that $f(x)$ remains the leftmost polynomial factor. Thus repeating the steps in Subsection 2.1 to $H'(n)$ yields again products of the form (27) with nesting depth $m-1$ with the following property: among all multiplicand representations of this new expression the $f(x)$ is still the leftmost polynomial.

Note further that also the new geometric product $G'(n)$ occurs with lower bound $\delta$. In order to turn it to a $1$-refined product, we may apply the transformations introduced in Section 2.2.

Finally, observe that in the special case $m = 1$, we get

\[
R(n) = \frac{H'(n)}{G'(n)} = \frac{g(n+1)}{g(\delta)} \in \mathbb{K}[n].
\]

Since the product $P'(n)$ itself might arise in the expression under consideration in the form $P'(n)^z$ with $z \in \mathbb{Z}$, we might introduce the factor $\frac{1}{P'(n)^z}$ in the final expression. However, since $R(n) \neq 0$ for all $n \geq \max(\delta - 1, 0)$, no extra poles will be introduced by this extra factor. Summarizing, also the final expression that has undergone the above transformation can be evaluated for all $n \geq \max(\delta - 1, 0)$. In particular, the input and output expression will have the same evaluation for each $n \geq \max(\delta - 1, 0)$.
More precisely we have the following:

\[ n \text{ holds for all } \]

\[ \in \hat{f} \text{ by Lemma } 1. \]

Using the relations (23), (22), (28), and (21) we can reduce all nesting depth 1 hypergeometric products whose multicaid representations are \( f_4(x), f_3(x), f_6(x), \) and \( f_2(x) \) respectively. More precisely we have the following:

\[ \prod_{k=3}^{n} f_4(k) = \prod_{k=3}^{n} (k + 2) = \frac{(n - 1)n(n + 1)(n + 2)}{24} \prod_{k=3}^{n} (k - 2) \] (29)

\[ \prod_{k=3}^{n} f_3(k) = \prod_{k=3}^{n} (k + 1) = \frac{(n - 1)n(n + 1)}{6} \prod_{k=3}^{n} (k - 2) \] (30)

\[ \prod_{k=3}^{n} f_6(k) = \prod_{k=3}^{n} k = \frac{(n - 1)n}{2} \prod_{k=3}^{n} (k - 2) \]

\[ \prod_{k=3}^{n} f_2(k) = \prod_{k=3}^{n} (k - 1) = (n - 1) \prod_{k=3}^{n} (k - 2). \]

Substituting (29), (30), (31), and (32) into (24), (25), and (26) and afterwards into the expression (19) gives

\[ \hat{H}(n) = \frac{2}{3} (n - 1)^3 n (n + 1)(n + 2) \prod_{k=1}^{n} 24^{-1/3} \prod_{k=1}^{n} 6^{-1/3} \left( \prod_{k=1}^{n} (k - 2) \right) \left( \prod_{k=3}^{n} (k + \frac{1}{2}) \right) \left( \prod_{k=3}^{n} (j - 2) \right). \] (33)

Note that \( H(n) = \hat{H}(n) \) for all \( n \geq 2 \). Furthermore, the distinct irreducible monic polynomials: \( (x - 2) \) and \( (x + \frac{1}{2}) \), that corresponds to the distinct innermost multiplicands of the products in factored form in \( \hat{H}(n) \) are shift-coprime among each other. Putting (17) and (18) in Example 4 and (33) together, we have that

\[ P(n) = \hat{P}(n) = \tilde{c} \tilde{r}(n) \hat{G}(n) \hat{H}(n) \] (34)

holds for all \( n \in \mathbb{Z}_{\geq 0} \) with \( n \geq 2 \), where the components of \( \hat{P}(n) \) are as follows:

\[ \tilde{c} = -\frac{254}{432}, \]

\[ \tilde{r}(n) = (n - 1)^3 n (n + 1)(n + 2), \]

\[ \hat{G}(n) = \left( \prod_{k=1}^{n} -1 \right) \left( \prod_{k=1}^{n} \sqrt{3} \right) \left( \prod_{k=1}^{n} 2 \right) \left( \prod_{k=1}^{n} 3 \right) \left( \prod_{k=1}^{n} 25 \right) \left( \prod_{k=1}^{n} k \right) \left( \prod_{k=1}^{n} \prod_{j=1}^{k} 5 \right) \left( \prod_{k=1}^{n} \prod_{j=1}^{k} 2 \right) \] (37)

\[ \hat{H}(n) = \left( \prod_{k=3}^{n} (k - 2) \right)^{3} \left( \prod_{k=3}^{n} (k + \frac{1}{2}) \right) \left( \prod_{k=3}^{n} k \right) \left( \prod_{k=3}^{n} (j - 2) \right). \] (38)

For further considerations the following definition will be convenient.

**Definition 5.** Let \( \mathbb{K}(x) \) be a rational function field and let \( H_1(n), \ldots, H_e(n) \in \text{ProdM}_n(\mathbb{K}(x)) \) where the arising hypergeometric products are in factored form. We say that \( H_1(n), \ldots, H_e(n) \) are in *shift-coprime product representation form* if
(1) the multiplicand representation of each product in $H_i(n)$ for $1 \leq i \leq e$ is an irreducible monic polynomial in $\mathbb{K}[x] \setminus \mathbb{K}$;

(2) the distinct multiplicand representations in $H_1(n), \ldots, H_e(n)$ are shift-coprime among each other.

Then the above symbolic manipulations can be summarized by the following method.

**Remark 2.** We are given $H_1(n), \ldots, H_e(n)$ in $\text{ProdM}_n(\mathbb{K}(x))$ where the products are in factored form and are all $\delta$-refined for some $\delta \in \mathbb{Z}_{\geq 0}$. In particular, suppose that all multiplicand representations are monic and that each $H_i(n)$ can be evaluated for $n \geq \nu$ for some $\nu \in \mathbb{Z}_{\geq 0}$ with $\nu \geq \delta$; note that such a $\nu$ can be derived by applying the available $\mathbb{Z}$-function to each rational function factor of $H_i(n)$ and taking its maximum value. Then one can follow the steps below to rewrite $H_1(n), \ldots, H_e(n)$ in a shift-coprime product representation form that yield the same evaluations for all $n \geq \nu$.

1. Factor the innermost multiplicand representations of all products in $H_1(n), \ldots, H_e(n)$ into irreducible monic polynomials in $\mathbb{K}[x] \setminus \mathbb{K}$ as in (6), and expand the product quantifier over the factorization as in (7). Let $\mathcal{D}$ be the set of all the irreducible monic polynomials.

2. Among all the irreducible monic polynomials in $\mathcal{D} \subseteq \mathbb{K}[x] \setminus \mathbb{K}$, compute the shift equivalence classes say, $\mathcal{E}_1, \ldots, \mathcal{E}_e$ with respect to the automorphism $\sigma(x) = x + 1$, and let $\mathcal{R}$ be the set of the leftmost polynomial of each equivalent class. Thus, the elements of the set $\mathcal{R}$ are shift-coprime among each other and each element represents exactly one equivalence class.

3. Among all the products over the irreducible monic polynomials obtained in step (1), take those with the highest nesting depth and reduce them with the elements in $\mathcal{R}$ following the construction of Remark [1]. During this rewriting one also obtains extra constants, geometric products and rational expressions from $\mathbb{K}(n)$ that are collected accordingly.

4. Go to step (1) and update the corresponding product expressions until the multiplicand representations of all products are in $\mathcal{R}$.

In particular, if $\mathbb{K}$ is computable and one can factorize polynomials over $\mathbb{K}$, all the above steps can be carried out explicitly.

Summarizing, given $\{P_1(n), \ldots, P_e(n)\} \subseteq \text{Prod}_n(\mathbb{K}(x))$, we can bring each $P_i(n)$ to the form (6) with $P'_i(n) := c_i, G_i(n), H_i(n)$ with $c_i \in \mathbb{K}^*$, $G_i(n) \in \text{ProdM}_n(\mathbb{K})$ and $H_i(n) \in \text{ProdM}_n(\mathbb{K}(x))$ such that $P_i(n) = P'_i(n)$ holds for all $n \geq \max(0, \delta - 1)$. Then applying the sketched algorithm in Remark 2 to $H_1(n), \ldots, H_e(n)$ we get the output $\tilde{H}_1(n), \ldots, \tilde{H}_e(n)$ where the $c_i$ and $\tilde{H}_i(n)$ are correspondingly updated to $\tilde{c}_i$ and $\tilde{H}_i(n)$ within Step 3 of Remark 2. In particular, we obtain for each component an extra factor $\tilde{r}_i(x) \in \mathbb{K}(x)$ where for $n \in \mathbb{Z}_{\geq 0}$ with $n \geq \max(0, \delta - 1)$, no poles are introduced in the evaluation of $\tilde{r}_i(n)$. Afterwards, another synchronization will be necessary to bring the new geometric products in $\tilde{H}_i(n)$ to 1-refined form (by again updating the $\tilde{c}_i$ accordingly). The final result can be summarized in the following proposition.

**Proposition 2.** Let $\mathbb{K}(x)$ be a rational function field and suppose that we are given the hypergeometric products $\{P_1(n), \ldots, P_e(n)\} \subseteq \text{Prod}_n(\mathbb{K}(x))$ of nesting depth at most $d \in \mathbb{Z}_{\geq 0}$ and there are

1. $\tilde{c}_1, \ldots, \tilde{c}_e \in \mathbb{K}^*$;

2. for all $1 \leq \ell \leq e$ rational functions $r_{\ell}(x) \in \mathbb{K}(x)^*$;

\[\text{As observed in Remark [1] the set } \mathcal{R} \text{ of the leftmost polynomials in step (2) does not change}\]
(3) geometric product expressions $\hat{G}_1(n), \ldots, \hat{G}_e(n) \in \text{ProdM}_n(\mathbb{K})$ which are all 1-refined;

(4) hypergeometric product expressions $\tilde{H}_1(n), \ldots, \tilde{H}_e(n) \in \text{ProdM}_n(\mathbb{K}(x))$ which are $\delta$-refined and are in shift-coprime product representation form.

such that for $1 \leq \ell \leq e$ and for all $n \geq \max(0, \delta - 1)$ we have

$$P_\ell(n) = \hat{c}_\ell \hat{v}_\ell(n) \tilde{G}_\ell(n) \tilde{H}_\ell(n) \neq 0. \quad (39)$$

If $\mathbb{K}$ is computable and one can factorize polynomials in $\mathbb{K}$, then $\delta$ and the above representation can be computed.

From now on, we assume that the arising hypergeometric products $P_1(n), \ldots, P_e(n) \in \text{Prod}_e(\mathbb{K}(x))$ have undergone the preprocessing steps discussed above yielding the representation given in (39). In general, there are still algebraic relations among the products that occur in the derived expressions (39) with $1 \leq \ell \leq e$, i.e., statements [2] and [3] of Problem RPE do not hold yet. In order to accomplish this task, extra insight from difference ring theory will be utilized. More precisely, we will show that the hypergeometric products coming from the $\tilde{H}_\ell$ are already algebraically independent, but the representation of the geometric products have to be improved to establish a solution of Problem RPE.

3. A naive difference ring approach: AP-extensions

Inspired by [Kar81, Sch16, Sch17], this Section focuses on an algebraic setting of difference rings (resp. fields) in which expressions of $\text{ProdM}_n(\mathbb{K}(x))$ can be naturally rephrased.

3.1. Difference Fields and Difference Rings

A difference ring (resp. field) $(\mathbb{A}, \sigma)$ is a ring (resp. field) $\mathbb{A}$ together with a ring (resp. field) automorphism $\sigma : \mathbb{A} \rightarrow \mathbb{A}$. Subsequently, all rings (resp. fields) are commutative with unity; in addition they contain the set of rational numbers $\mathbb{Q}$, as a subring (resp. subfield). The multiplicative group of units of a ring (resp. field) $\mathbb{A}$ is denoted by $\mathbb{A}^*$. A difference ring (resp. field) $(\mathbb{A}, \sigma)$ is called computable if $\mathbb{A}$ and $\sigma$ are both computable. In the following we will introduce AP-extensions that will be the foundation to represent hypergeometric products of finite nesting depth in difference rings.

A-extensions will be used to cover objects like $\zeta^k$ where $\zeta$ is a root of unity. In general, let $(\mathbb{A}, \sigma)$ be a difference ring and let $\zeta \in \mathbb{A}^*$ be a $\lambda$-th root of unity with $\lambda > 1$ (i.e., $\lambda \in \mathbb{Z}_{\geq 2}$ with $\zeta^\lambda = 1$). Take the uniquely determined difference ring extension $(\mathbb{A}[y], \sigma)$ of $(\mathbb{A}, \sigma)$ where $y$ is transcendental over $\mathbb{A}$ and $\sigma(y) = \zeta y$. Now consider the ideal $I := \langle y^\lambda - 1 \rangle$ and the quotient ring $\mathbb{E} := \mathbb{A}[y]/I$. Since $I$ is closed under $\sigma$ and $\sigma^{-1}$ i.e., $I$ is a reflexive difference ideal, we can define the map $\sigma : \mathbb{E} \rightarrow \mathbb{E}$ with $\sigma(h + I) = \sigma(h) + I$ which forms a ring automorphism. Note that by this construction the ring $\mathbb{A}$ can naturally be embedded into the ring $\mathbb{E}$ by identifying $a \in \mathbb{A}$ with $a + I \in \mathbb{E}$, i.e., $a \mapsto a + I$. Now set $\vartheta := y + I$. Then $(\mathbb{A}[\vartheta], \sigma)$ is a difference ring extension of $(\mathbb{A}, \sigma)$ subject to the relations $\vartheta^\lambda = 1$ and $\sigma(\vartheta) = \zeta \vartheta$. This extension is called an algebraic extension (in short A-extension) of order $\lambda$. The generator $\vartheta$ is called an A-monomial with its order $\lambda = \min\{n > 0 | \vartheta^n = 1\}$. Note that the ring $\mathbb{A}[\vartheta]$ is not an integral domain (i.e., it has zero-divisors) since $(\vartheta - 1)(\vartheta^{\lambda-1} + \cdots + \vartheta + 1) = 0$ but $(\vartheta - 1) \neq 0 \neq (\vartheta^{\lambda-1} + \cdots + \vartheta + 1)$. In this setting, the A-monomial $\vartheta$ with the relations $\vartheta^\lambda = 1$ and $\sigma(\vartheta) = \zeta \vartheta$ with $\zeta := e^{\frac{2\pi i}{\lambda}} = (-1)^\frac{\lambda}{2}$, models $\zeta^k$ subject to the relations $(\zeta^k)^\lambda = 1$ and $\zeta^{k+1} = \zeta^{k} \zeta^k$.

In addition, we define P-extensions in order to treat products of finite nesting depth whose multiplicands are not given by roots of unity. Let $(\mathbb{A}, \sigma)$ be a difference ring, $\alpha \in \mathbb{A}^*$ be a unit, and consider the ring of Laurent polynomials $\mathbb{A}[t, t^{-1}]$ (i.e., $t$ is transcendental over $\mathbb{A}$). Then there is a unique difference ring extension $(\mathbb{A}[t, t^{-1}], \sigma)$ of $(\mathbb{A}, \sigma)$ with $\sigma(t) = \alpha t$ and $\sigma(t^{-1}) = \alpha^{-1} t^{-1}$. The extension here is called a product-extension (in short P-extension) and the generator $t$ is called a P-monomial.

We introduce the following notations for convenience. Let $(\mathbb{E}, \sigma)$ be a difference ring extension of $(\mathbb{A}, \sigma)$ with $t \in \mathbb{E}$. $\mathbb{A}(t)$ denotes the ring of Laurent polynomials $\mathbb{A}[t, \frac{1}{t}]$ (i.e., $t$ is transcendental over...
\( \mathbb{A} \) if \( \mathbb{A}[t, \frac{1}{t} \mid \sigma] \) is a \( \mathbb{P} \)-extension of \((\mathbb{A}, \sigma)\). Lastly, \( \mathbb{A}(t) \) denotes the ring \( \mathbb{A}[t] \) with \( t \notin \mathbb{A} \) but subject to the relation \( t^\lambda = 1 \) if \((\mathbb{A}[t], \sigma)\) is an \( \mathbb{A} \)-extension of \((\mathbb{A}, \sigma)\) of order \( \lambda \).

We say that the difference ring extension \((\mathbb{A}(t), \sigma)\) of \((\mathbb{A}, \sigma)\) is an \( \mathbb{AP} \)-extension (and \( t \) is an \( \mathbb{AP} \)-monomial) if it is an \( \mathbb{A} \)- or a \( \mathbb{P} \)-extension. Finally, we call \((\mathbb{A}(t_1) \ldots \langle t_e \rangle, \sigma)\) a (nested) \( \mathbb{A}/\mathbb{P}/\mathbb{AP} \)-extension of \((\mathbb{A}, \sigma)\) if it is built by a tower of such extensions.

In the following we will restrict to the subclass of ordered simple \( \mathbb{AP} \)-extension. Here, the following definitions are useful.

**Definition 6.** Let \((E, \sigma)\) be a (nested) \( \mathbb{AP} \)-extension of \((\mathbb{A}, \sigma)\) with \( E = \mathbb{A}(t_1) \ldots \langle t_e \rangle \) where \( \sigma(t_i) = \alpha_i t_i \) for \( 1 \leq i \leq e \). We define the **depth** function of elements of \( E \) over \( \mathbb{A} \), \( \delta : E \to \mathbb{Z}_{\geq 0} \) as follows:

1. For any \( h \in \mathbb{A} \), \( \delta_h(h) = 0 \).
2. If \( \delta_h \) is defined for \((\mathbb{A}(t_1) \ldots \langle t_{i-1} \rangle, \sigma)\) with \( i > 1 \), then we define \( \delta_h(t_i) := \delta_h(\alpha_i) + 1 \) and for \( f \in \mathbb{A}(t_1) \ldots \langle t_{i-1} \rangle \), we define \( \delta_h(f) := \max \{ \{ \delta_h(t_i) \mid t_i \text{ occurs in } f \} \cup \{0\} \} \).

The **extension depth** of \((E, \sigma)\) over \( \mathbb{A} \) is given by \( \delta_h(E) := (\delta_h(t_1), \ldots, \delta_h(t_e)) \). We call such an extension ordered, if \( \delta_h(t_1) \leq \delta_h(t_2) \leq \cdots \leq \delta_h(t_e) \). In particular, we say that \((E, \sigma)\) is of **monomial depth** \( m \) if \( m = \max(0, \delta_h(t_1), \ldots, \delta_h(t_e)) \). If \( \mathbb{A} \) is clear from the context, we write \( \delta_h \) as \( \delta \).

Now, let \((E, \sigma)\) with \( E = \mathbb{A}(t_1) \ldots \langle t_e \rangle \) be a nested \( \mathbb{A}/\mathbb{P}/\mathbb{AP} \)-extension of a difference ring \((\mathbb{A}, \sigma)\) and let \( G \) be a multiplicative subgroup of \( \mathbb{A}^* \). Following \[Sch16, Sch17\] we call

\[
G_E^\mathbb{A} := \{ g t_1^{v_1} \cdots t_e^{v_e} \mid g \in G, \text{ and } v_i \in \mathbb{Z} \}
\tag{40}
\]

the **product group** over \( G \) with respect to \( \mathbb{A}/\mathbb{P}/\mathbb{AP} \)-monomials for the nested \( \mathbb{A}/\mathbb{P}/\mathbb{AP} \)-extension \((E, \sigma)\) of \((\mathbb{A}, \sigma)\). In the following we will restrict ourselves to the following subclass of \( \mathbb{AP} \)-extensions.

**Definition 7.** Let \((\mathbb{A}, \sigma)\) be a difference ring and let \( G \) be a subgroup of \( \mathbb{A}^* \). Let \((E, \sigma)\) be an \( \mathbb{A}/\mathbb{P}/\mathbb{AP} \)-extension of \((\mathbb{A}, \sigma)\) with \( E = \mathbb{A}(t_1) \ldots \langle t_e \rangle \). Then this extension is called \( G \)-**simple** if for all \( 1 \leq i \leq e \),

\[
\frac{\sigma(t_i)}{t_i} \in G_{\mathbb{A}(t_1) \ldots \langle t_{i-1} \rangle}^\mathbb{A}.
\]

In addition such a \( G \)-simple extension is called \( G \)-**basic**\(^4\) if for any \( \mathbb{A} \)-monomial \( t_i \) we have \( \frac{\sigma(t_i)}{t_i} \in \text{const}(\mathbb{A}, \sigma)^* \) and for any \( \mathbb{P} \)-monomial \( t_i \) we have that \( \frac{\sigma(t_i)}{t_i} \in G_{\mathbb{A}(t_1) \ldots \langle t_{i-1} \rangle}^\mathbb{A} \) is free of \( \mathbb{A} \)-monomials. If \( G = \mathbb{A}^* \), such extensions are also called **simple** (resp. **basic**) instead of \( \mathbb{A}^* \)-**simple** (\( \mathbb{A}^* \)-**basic**).

In particular, we will work with the following class of simple \( \mathbb{AP} \)-extensions that are closely related to the products in factored form given in \[Z\]; for concrete constructions see Example \[Z\] below.

**Definition 8.** Let \((\mathbb{A}, \sigma)\) be a difference ring and \( G \) be a subgroup of \( \mathbb{A}^* \). We call \((\mathbb{A}(t_1) \ldots \langle t_e \rangle, \sigma)\) a **single chain** \( \mathbb{A}/\mathbb{P}/\mathbb{AP} \)-extension of \((\mathbb{A}, \sigma)\) over \( G \) if for all \( 1 \leq k \leq e \),

\[
\sigma(t_k) = c_k t_1 \cdots t_{k-1} t_k, \quad \text{with } c_k \in G.
\]

We call \( c_1 \) also the **base** of the single chain \( \mathbb{A}/\mathbb{P}/\mathbb{AP} \)-extension. If \( G = \mathbb{A}^* \), we also say that \((\mathbb{A}(t_1) \ldots \langle t_e \rangle, \sigma)\) is a **single chain** \( \mathbb{A}/\mathbb{P}/\mathbb{AP} \)-extension of \((\mathbb{A}, \sigma)\). Further, we call \((E, \sigma)\) a **multiple chain** \( \mathbb{A}/\mathbb{P}/\mathbb{AP} \)-extension of \((\mathbb{A}, \sigma)\) over \( G \) with base \( (c_1, \ldots, c_m) \in G^m \) if it is a tower of \( m \) single chain \( \mathbb{A}/\mathbb{P} \)-extensions over \( G \) with the bases \( c_1, \ldots, c_m \), respectively. If \( G = \mathbb{A}^* \), we simply call it a multiple chain \( \mathbb{A}/\mathbb{P}/\mathbb{AP} \)-extension.

\(^4\)In other words, products whose multiplicands are roots of unity have nesting depth 1 (and the roots of unity are from the constant field), while the remaining products do not depend on these products over roots of unity.
Remark 3. Let $(\mathbb{A}, (t_1, \ldots, t_e), \sigma)$ be a single chain $A$-P/A-AP-extension of $(\mathbb{A}, \sigma)$ as given in Definition 8 and let $\delta : \mathbb{A} \langle t_1, \ldots, t_e \rangle \to \mathbb{Z}_{\geq 0}$ be the depth function over $\mathbb{A}$. Then we have $\delta(t_k) = k$ for all $1 \leq k \leq e$. In particular, the extension is ordered, its extension-depth is $(1, 2, \ldots, e)$ and the monomial depth is $e$. Furthermore observe that for $2 \leq i \leq e$ we have

$$
\sigma(t_i) = \sigma(t_{i-1}) t_i \Leftrightarrow \frac{t_i}{\sigma^{-1}(t_i)} = t_{i-1}.
$$

3.2. Ring of Sequences

For a field $\mathbb{K}$ we denote by $\mathbb{K}^{\mathbb{Z}_{\geq 0}}$ the set of all sequences

$$
\langle a(n) \rangle_{n \geq 0} = \langle a(0), a(1), a(2), \ldots \rangle \tag{41}
$$

whose terms are in $\mathbb{K}$. Equipping $\mathbb{K}^{\mathbb{Z}_{\geq 0}}$ with component-wise addition and multiplication, we get a commutative ring. In this ring, the field $\mathbb{K}$ can be naturally embedded into $\mathbb{K}^{\mathbb{Z}_{\geq 0}}$ as a subring, by identifying any $c \in \mathbb{K}$ with the constant sequence $(c, c, c, \ldots) \in \mathbb{K}^{\mathbb{Z}_{\geq 0}}$. Following the construction in [PWZ96 Section 8.2], we turn the shift operator $S : \mathbb{K}^{\mathbb{Z}_{\geq 0}} \to \mathbb{K}^{\mathbb{Z}_{\geq 0}}$ with

$$
S : \langle a(0), a(1), a(2), \ldots \rangle \mapsto \langle a(1), a(2), a(3), \ldots \rangle \tag{42}
$$

into a ring automorphism by introducing an equivalence relation $\sim$ on sequences in $\mathbb{K}^{\mathbb{Z}_{\geq 0}}$. Two sequences $A = \langle a(n) \rangle_{n \geq 0}$ and $B = \langle b(n) \rangle_{n \geq 0}$ are said to be equivalent (in short $A \sim B$) if and only if there exists a non-negative integer $\delta$ such that

$$
\forall n \geq \delta : a(n) = b(n).
$$

The set of equivalence classes form a ring again with component-wise addition and multiplication which will we denote by $\mathcal{S}(\mathbb{K}) := \mathbb{K}^{\mathbb{Z}_{\geq 0}}/\sim$. Now it is obvious that $S : \mathcal{S}(\mathbb{K}) \to \mathcal{S}(\mathbb{K})$ with $\langle 42 \rangle$ is bijective and thus a ring automorphism. We call $(\mathcal{S}(\mathbb{K}), S)$ also the difference ring of sequences over $\mathbb{K}$. For simplicity, we denote the elements of $\mathcal{S}(\mathbb{K})$ by the usual sequence notation as in $\langle 41 \rangle$ above.

We will follow the convention introduced in [PST19] to illustrate how the indefinite products of finite nesting depth covered in this article are modelled by expressions in a difference ring.

Definition 9. Let $(\mathbb{A}, \sigma)$ be a difference ring with a constant field $\mathbb{K} = \text{const}(\mathbb{A}, \sigma)$. An evaluation function $ev : \mathbb{A} \times \mathbb{Z}_{\geq 0} \to \mathbb{K}$ for $(\mathbb{A}, \sigma)$ is a function which satisfies the following three properties:

(i) for all $c \in \mathbb{K}$, there is a natural number $\delta \geq 0$ such that

$$
\forall n \geq \delta : ev(c, n) = c; \tag{43}
$$

(ii) for all $f, g \in \mathbb{A}$ there is a natural number $\delta \geq 0$ such that

$$
\forall n \geq \delta : ev(f g, n) = ev(f, n) ev(g, n),

\forall n \geq \delta : ev(f + g, n) = ev(f, n) + ev(g, n); \tag{44}
$$

(iii) for all $f \in \mathbb{A}$ and $i \in \mathbb{Z}$, there is a natural number $\delta \geq 0$ such that

$$
\forall n \geq \delta : ev(\sigma^i(f), n) = ev(f, n + i). \tag{45}
$$

We say a sequence $\langle F(n) \rangle_{n \geq 0} \in \mathcal{S}(\mathbb{K})$ is modelled by $f \in \mathbb{A}$ in the difference ring $(\mathbb{A}, \sigma)$, if there is an evaluation function $ev$ such that

$$
F(k) = ev(f, k)
$$

holds for all $k \in \mathbb{Z}_{\geq 0}$ from a certain point on.
In this article, our base field is a rational function field $\mathbb{K}(x)$ which is equipped with the evaluation function $ev : \mathbb{K}(x) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}$ defined as follows. For $f = \frac{g}{h} \in \mathbb{K}(x)$ with $h \neq 0$ where $g$ and $h$ are coprime (if $g = 0$ we take $h = 1$) we have

$$ev(f, k) := \begin{cases} 0 & \text{if } h(k) = 0, \\ \frac{g(k)}{h(k)} & \text{if } h(k) \neq 0. \end{cases}$$ (46)

Here, $g(k)$ and $h(k)$ are the usual polynomial evaluation at some natural number $k$.

Then given a tower of AP-extension defined over $(\mathbb{K}(x), \sigma)$, one can define an appropriate evaluation function by iterative applications of the following lemma that is implied by [Sch17, Lemma 5.4].

**Lemma 2.** Let $(\mathbb{A}, \sigma)$ be a difference ring with constant field $\mathbb{K}$ and let $ev : \mathbb{A} \times \mathbb{Z}_{\geq 0} \to \mathbb{K}$ be an evaluation function for $(\mathbb{A}, \sigma)$. Let $(\mathbb{A}(t), \sigma)$ be an AP-extension of $(\mathbb{A}, \sigma)$ with $\sigma(t) = \alpha t$ ($\alpha \in \mathbb{A}^*$) and suppose that there is a $\delta \in \mathbb{Z}_{\geq 0}$ such that $ev(\alpha, n) \neq 0$ for all $n \geq \delta$. Further, take $u \in \mathbb{K}^*$; if $u^\lambda = 1$ for some $\lambda > 1$, we further assume that $u^\lambda = 1$ holds. Consider the map $ev' : \mathbb{A}(t) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}$ defined by

$$ev'(\sum_i h_i t^i, n) = \sum_i ev(h_i, n) ev'(t, n)^i$$

with $ev'(t, n) = u \prod_{k=\delta}^n ev(\alpha, k - 1)$. Then $ev'$ is an evaluation function for $(\mathbb{A}(t), \sigma)$.

We summarize the above constructions with the following example.

**Example 7.** Let $\mathbb{K} = \mathbb{Q}(\sqrt{3})$ and take the difference field $(\mathbb{K}(x), \sigma)$ where the automorphism is defined by $\sigma(x) = x + 1$ and $\sigma|_\mathbb{K} = \text{id}$. Furthermore, take the evaluation function $ev : \mathbb{K}(x) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}$ given by (46). Then we can construct the following single chain extensions of $(\mathbb{K}(x), \sigma)$ in order to model the geometric product $\hat{G}(n)$ and the hypergeometric product $\hat{H}(n)$ given in (37) and (38).

1. We define the single chain A-extension $(\mathbb{K}(x)(\vartheta_{1,1}), \vartheta_{1,2}, \sigma)$ of $(\mathbb{K}(x), \sigma)$ over $\mathbb{K}$ of order 2, based at $-1$ where the automorphism is given by

   $$\sigma(\vartheta_{1,1}) = -\vartheta_{1,1}, \quad \sigma(\vartheta_{1,2}) = -\vartheta_{1,1} \vartheta_{1,2}. \quad (47)$$

   In addition by applying Lemma 2 twice we extend the evaluation function to $ev : \mathbb{K}(x)(\vartheta_{1,1}, \vartheta_{1,2}) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}$ with

   $$ev(\vartheta_{1,1}, n) = \prod_{k=1}^n (-1), \quad ev(\vartheta_{1,2}, n) = \prod_{k=1}^n (-1). \quad (48)$$

2. Similarly, define the single chain P-extension $(\mathbb{K}(x)(y_{1,1}), \sigma)$ of $(\mathbb{K}(x), \sigma)$ over $\mathbb{K}$ based at $\sqrt{3}$ together with the evaluation function $ev : \mathbb{K}(x)(y_{1,1}) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}$ (using Lemma 2) by

   $$\sigma(y_{1,1}) = \sqrt{3} y_{1,1}, \quad \text{and} \quad ev(y_{1,1}, n) = \prod_{k=1}^n \sqrt{3}. \quad (49)$$

3. Define the single chain P-extension $(\mathbb{K}(x)(y_{2,1}, y_{2,2}), \sigma)$ of $(\mathbb{K}(x), \sigma)$ over $\mathbb{K}$ based at 2 equipped with the evaluation function $ev : \mathbb{K}(x)(y_{2,1}, y_{2,2}) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}$ by

   $$\sigma(y_{2,1}) = 2 y_{2,1}, \quad \sigma(y_{2,2}) = 2 y_{2,1} y_{2,2}, \quad \text{and} \quad ev(y_{2,1}, n) = \prod_{k=1}^n 2, \quad ev(y_{2,2}, n) = \prod_{k=1}^n 2. \quad (50)$$
(4) Define the single chain P-extension \((\mathbb{K}(x)(y_{3,1}), \sigma)\) of \((\mathbb{K}(x), \sigma)\) over \(\mathbb{K}\) based at 3 together with the evaluation function \(ev : \mathbb{K}(x)(y_{3,1}) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}\) by

\[
\sigma(y_{3,1}) = 3y_{3,1}, \quad \text{and} \quad ev(y_{3,1}, n) = \prod_{k=1}^{n} 3.
\]

(5) Define the single chain P-extension \((\mathbb{K}(x)(y_{4,1}), \sigma)\) of \((\mathbb{K}(x), \sigma)\) over \(\mathbb{K}\) with 5 as its base accompanied with the evaluation function \(ev : \mathbb{K}(x)(y_{4,1}) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}\) by

\[
\sigma(y_{4,1}) = 5y_{4,1}, \quad \sigma(y_{4,2}) = 5y_{4,1}y_{4,2}, \quad \text{and} \quad ev(y_{4,1}, n) = \prod_{k=1}^{n} 5, \quad ev(y_{4,1}, n) = \prod_{k=1}^{n} \prod_{j=1}^{k} 5.
\]

(6) Define the single chain P-extension \((\mathbb{K}(x)(y_{5,1}), \sigma)\) of \((\mathbb{K}(x), \sigma)\) over \(\mathbb{K}\) with 25 as its base together with the evaluation function \(ev : \mathbb{K}(x)(y_{5,1}) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}\) by

\[
\sigma(y_{5,1}) = 25y_{5,1}, \quad \text{and} \quad ev(y_{5,1}, n) = \prod_{k=1}^{n} 25.
\]

(7) Define the single chain P-extension \((\mathbb{K}(x)(z_{1,1}), \sigma)\) of \((\mathbb{K}(x), \sigma)\) over \(\mathbb{K}(x)\) based at \((x - 2)\) and the evaluation function \(ev : \mathbb{K}(x)(z_{1,1}) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}\) by

\[
\sigma(z_{1,1}) = (x - 1)z_{1,1}, \quad \sigma(z_{1,2}) = (x - 1)z_{1,1}z_{1,2}, \quad \text{and} \quad ev(z_{1,1}, n) = \prod_{k=3}^{n} (k - 2), \quad ev(z_{1,2}, n) = \prod_{k=3}^{n} \prod_{j=3}^{k} (j - 2).
\]

(8) Define the single chain P-extension \((\mathbb{K}(x)(z_{2,1}), \sigma)\) of \((\mathbb{K}(x), \sigma)\) over \(\mathbb{K}(x)\) with \((x + \frac{25}{24})\) as its base together with the evaluation function \(ev : \mathbb{K}(x)(z_{2,1}) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}\) by

\[
\sigma(z_{2,1}) = (x + \frac{25}{24})z_{2,1}, \quad \text{and} \quad ev(z_{1,1}, n) = \prod_{k=3}^{n} (k + \frac{1}{24}).
\]

Putting everything together, we have constructed the multiple chain AP-extension \((A, \sigma)\) of \((\mathbb{K}(x), \sigma)\) with

\[
A = \mathbb{K}(x) \langle \partial_{1,1}, \partial_{1,2} \rangle \langle y_{1,1}, y_{2,1} \rangle \langle y_{1,2}, y_{2,2} \rangle \langle y_{3,1}, y_{3,2}, y_{4,1}, y_{4,2} \rangle \langle z_{1,1}, z_{1,2}, z_{2,1}, z_{2,2} \rangle
\]

based at \((-1, \sqrt{3}, 2, 2, 3, 5, 5, 25, x - 2, x - 2, x + \frac{1}{24}\) where \((1, 2, 1, 1, 2, 1, 1, 2, 1, 1, 2, 1, 1)\) is the extension depth. In this ring, the geometric product \(G(n)\) and the hypergeometric product \(H(n)\) defined in \((37)\) and \((38)\) are modelled by

\[
g = \frac{\partial_{1,1} y_{3,1} y_{3,1} \partial_{1,2} y_{2,2}}{y_{1,1} y_{2,1} y_{4,2}} \quad \text{and} \quad h = z_{1,1}^3 z_{2,1} z_{1,2}
\]

respectively. That is, \(G(n) = ev(g, n)\) holds for all \(n \geq 1\) and \(H(n) = ev(h, n)\) holds for all \(n \geq 2\). As a consequence, the indefinite hypergeometric product expression \(\tilde{P}(n)\) defined in \((54)\) is modelled by the expression

\[
\tilde{p} = \frac{254 (x - 1)^3 x (x + 1) (x + 2)}{432} \quad gh \in A.
\]

This means that \(\tilde{P}(n) = ev(\tilde{p}, n)\) holds for all \(n \in \mathbb{Z}_{\geq 0}\) with \(n \geq 2\).
In this way, all products arising in the rewritten P-extension within the constructed difference ring (A-P-extensions further to the class of RΠ-extensions. In this regard, the set of constants algebraic relations among the monomials. In order to tackle Problem RPE above, we will refine if const(A) = \{c \in A \mid \sigma(e) = c\}

of a difference ring (field) (A, \sigma) plays a decisive role. In general it forms a subring (subfield) of A which contains the rational numbers \(\mathbb{Q}\) as subfield. In this article, const(A, \sigma) will always be a field also called the constant field of (A, \sigma), which we will also denote by \(\mathbb{K}\). We note further that one can decide if \(c \in A\) is a constant if (A, \sigma) is computable.

We are now ready to refine AP-extensions as follows.

Definition 10. Let \((A, \sigma)\) be an A-/P-/AP-extension of \((A, \sigma)\). Then it is called an R-/Π-/RII-extension if const(A, \sigma) = const(A, \sigma). Depending on the type of extension, we call the generator t an R-/Π-/RII-monomial, respectively. A (nested) A-/P-/AP-extension \((A, \langle t_1 \rangle, \ldots, (t_e), \sigma)\) of \((A, \sigma)\) with const(A, \langle t_1 \rangle, \ldots, (t_e), \sigma) = const(A, \sigma) is also called a (nested) R-/Π-/RII-extension.

Remark 4. We can carry out the following construction to model the given hypergeometric products \(P_1(n), \ldots, P_e(n) \in \text{Prod}_n(\mathbb{K}(x))\) of finite nesting depth and an expression in terms of these products in a P-extension.\(^5\) Here we start with the difference field \((\mathbb{K}(x), \sigma)\) given by \(\sigma|_{\mathbb{K}} = \text{id}\) and \(\sigma(x) = x + 1\) which is equipped with the evaluation function \(ev\) given in \([10]\) and the zero-function \([9]\).

1. For \(1 \leq i \leq e\), rewrite \(P_i(n)\) such that it is composed multiplicatively by products in factored form and such that all products are \(\delta\)-refined (in its strongest form, one may use the representation given in Proposition 2). Let \(P\) be the set of all products that occur in the rewritten \(P_i(n)\).

2. Among all nested products of \(P\) with the same multiplicand representation \(c(x) \in \mathbb{K}(x)\), take one of the products, say \(F(n)\), with the highest nesting depth \(m\). Construct the corresponding single chain P-extension \((\mathbb{K}(x), \langle p_1 \rangle, \ldots, \langle p_m \rangle), \sigma)\) of \((\mathbb{K}(x), \sigma)\) over \(\mathbb{K}(x)\) and extend the evaluation function accordingly such that the outermost P-monomial \(p_m\) models \(F(n)\) for all \(n \geq \delta\). In particular, any arising product with the same multiplicand representation and depth \(\mu < m\) is modelled by \(p_\mu\). Thus we can remove all the products from \(P\) which have the same multiplicand \(c(x)\).

3. Repeat step [2] for the remaining elements of \(P\).

4. Combine the constructed single chain P-extensions of \((\mathbb{K}(x), \sigma)\) over \(\mathbb{K}(x)\) to obtain a multiple chain P-extension \((A, \sigma)\) of \((\mathbb{K}(x), \sigma)\) over \(\mathbb{K}(x)\). In addition, combine the evaluation functions to one extended version.

In this way, all products arising in the rewritten \(P_1(n), \ldots, P_e(n)\) can be modelled by a P-monomial within the constructed difference ring \((A, \sigma)\). Furthermore, consider any expression \(A(n)\) of the form \([2]\) in terms of the original products \(P_1(n), \ldots, P_e(n)\) and let \(\lambda = \max\{|Z(a_v)\mid v \in S\}\), i.e., the evaluation \(a_v(n)\) does not introduce poles for any \(n \in \mathbb{Z}_{\geq 0}\) with \(n \geq \lambda\). Then replacing the rewritten products in \(A(n)\) and afterwards replacing the involved products by the corresponding P-monomials yields a given \(a \in A\) with \(ev(a, n) = A(n)\) for all \(n \geq \max(\delta, \lambda)\). In particular, if \(\mathbb{K}\) is computable and the zero-function \(Z\) is computable, this construction can be given explicitly.

4. A refined difference ring approach: RΠ-extensions

In general, the naive construction of an (ordered) multiple chain P-extension \((A, \sigma)\) of \((\mathbb{K}(x), \sigma)\) following Remark 4 or a slightly refined construction of an AP-extension like in Example 7 introduce algebraic relations among the monomials. In order to tackle Problem RPE above, we will refine AP-extensions further to the class of RΠ-extensions. In this regard, the set of constants

\[\text{const}(A, \sigma) = \{c \in A \mid \sigma(c) = c\}\]

\(^5\)Similarly, one can carry out such a construction for A-extensions or AP-extensions by checking in addition if the arising multiplicands are built by roots of unity.
Given an A-/P-/AP-extension, there exist criteria that enable one able to check if it is an R-/Π-/RII-extension. We refer the reader to see [Sch16 Theorem 2.12] for further details and proofs.

**Theorem 1.** Let \((\mathcal{A}, \sigma)\) be a difference ring. Then the following statements hold.

1. Let \((\mathcal{A}[t], \sigma)\) be a P-extension of \((\mathcal{A}, \sigma)\) with \(\sigma(t) = \alpha t\) where \(\alpha \in \mathcal{A}^*\). Then this is a Π-extension (i.e., \(\text{const}(\mathcal{A}[t], \sigma) = \text{const}(\mathcal{A}, \sigma)\)) iff there are no \(g \in \mathcal{A} \setminus \{0\}\) and \(v \in \mathbb{Z} \setminus \{0\}\) with \(\sigma(g) = \alpha^v g\).

2. Let \((\mathcal{A}[\vartheta], \sigma)\) be an A-extension of \((\mathcal{A}, \sigma)\) of order \(\lambda > 1\) with \(\sigma(\vartheta) = \zeta \vartheta\) where \(\zeta \in \mathcal{A}^*\). Then this is an R-extension (i.e., \(\text{const}(\mathcal{A}[\vartheta], \sigma) = \text{const}(\mathcal{A}, \sigma)\)) iff there are no \(g \in \mathcal{A} \setminus \{0\}\) and \(v \in \{1, \ldots, \lambda - 1\}\) with \(\sigma(g) = \zeta^v g\). If it is an R-extension, \(\zeta\) is a primitive \(\lambda\)-th root of unity.

We remark that the above definitions and also Theorem 1 are inspired by Karr’s ΠΣ-field extensions [Kar81, Sch01]. Since we will use this notion later (see Theorem 4 and 5 below) we will introduce them already here.

**Definition 11.** Let \((\mathcal{F}(t), \sigma)\) be a difference field extension of a difference field \((\mathcal{F}, \sigma)\) with \(t\) transcendental over \(\mathcal{F}\) and \(\sigma(t) = \alpha t + \beta\) with \(\alpha \in \mathbb{F}^*\) and \(\beta \in \mathbb{F}\). This extension is called a Σ-field extension if \(\alpha = 1\) and \(\text{const}(\mathcal{F}(t), \sigma) = \text{const}(\mathcal{F}, \sigma)\), and it is called a Π-field extension if \(\beta = 0\) and \(\text{const}(\mathcal{F}(t), \sigma) = \text{const}(\mathcal{F}, \sigma)\). A difference field \((\mathbb{K}(t_1), \ldots, t_e), \sigma\) is called a ΠΣ-field over \(\mathbb{K}\) if \((\mathbb{K}(t), \ldots, t), \sigma\) is a ΠΣ-extension of \((\mathbb{K}(t_1), \ldots, t_{i-1}), \sigma\) for \(1 \leq i \leq e\) with \(\text{const}(\mathbb{K}, \sigma) = \mathbb{K}\).

Throughout this article, our base case difference field is \((\mathbb{K}(x), \sigma)\) with the automorphism \(\sigma(x) = x + 1\) and \(\sigma|_{\mathbb{K}} = \text{id}\) which in fact is a Σ-extension of \((\mathbb{K}, \sigma)\), i.e., \(\text{const}(\mathbb{K}(x), \sigma) = \mathbb{K}\). In particular \((\mathbb{K}(x), \sigma)\) is a ΠΣ-field over \(\mathbb{K}\). We conclude this subsection by observing that the check if an A-extension is an R-extension (see part 2 of Theorem 1) is not necessary if the ground field is a ΠΣ-field; compare [OS18, Lemma 2.1].

**Lemma 3.** Let \((\mathcal{F}, \sigma)\) be a ΠΣ-field over \(\mathbb{K}\). Then any A-extension \((\mathcal{F}[\vartheta], \sigma)\) of \((\mathcal{F}, \sigma)\) with order \(\lambda > 1\) is an R-extension.

### 4.1. Embedding into the ring of sequences

In this subsection, we will discuss the connection between RII-extensions and the difference ring of sequences. More precisely, we will show how RII-extensions can be embedded into the difference ring of sequences [Sch17, compare also [PS97]. This feature will enable us to handle condition (2) of Problem RPE in the sections below.

**Definition 12.** Let \((\mathcal{A}, \sigma)\) and \((\mathcal{A}', \sigma')\) be two difference rings. The map \(\tau: \mathcal{A} \to \mathcal{A}'\) is called a difference ring homomorphism if \(\tau\) is a ring homomorphism, and for all \(f \in \mathcal{A}\), \(\tau(\sigma(f)) = \sigma'(\tau(f))\). If \(\tau\) is injective, then it is called a difference ring monomorphism or a difference ring embedding\(^6\). If \(\tau\) is a bijection, then it is a difference ring isomorphism and we say \((\mathcal{A}, \sigma)\) and \((\mathcal{A}', \sigma')\) are isomorphic; we write \((\mathcal{A}, \sigma) \simeq (\mathcal{A}', \sigma')\). Let \((\mathcal{E}, \sigma)\) and \((\mathcal{E}, \tilde{\sigma})\) be difference ring extensions of \((\mathcal{A}, \sigma)\). Then a difference ring-homomorphism/isomorphism/monomorphism \(\tau: \mathcal{E} \to \mathcal{E}\) is called an A-homomorphism/A-isomorphism/A-monomorphism, if \(\tau|_{\mathcal{A}} = \text{id}\).

Let \((\mathcal{A}, \sigma)\) be a difference ring with constant field \(\mathbb{K}\). A difference ring homomorphism (resp. monomorphism) \(\tau: \mathcal{A} \to S(\mathbb{K})\) is called \(\mathbb{K}\)-homomorphism (resp. -monomorphism) if for all \(c \in \mathbb{K}\) we have that \(\tau(c) = c := \langle c, c, c, \ldots \rangle\).

The following lemmas provide the key property that will enable us to embed RII-extensions into the ring of sequences. First, we recall that the evaluation function of a difference ring establishes naturally a \(\mathbb{K}\)-homomorphism. More precisely, by [Sch01 Lemma 2.5.1] we get

\(^6\)In this case, \((\tau(\mathcal{A}), \sigma)\) is a sub-difference ring of \((\mathcal{A}', \sigma')\) where \((\mathcal{A}, \sigma)\) and \((\tau(\mathcal{A}), \sigma)\) are the same up to renaming with respect to \(\tau\).
Lemma 4. Let \((A, \sigma)\) be a difference ring with constant field \(K\). Then the map \(\tau : A \rightarrow S(K)\) is a \(K\)-homomorphism if and only if there is an evaluation function \(ev : A \times \mathbb{Z}_{\geq 0} \rightarrow K\) for \((A, \sigma)\) (see Definition 2) with \(\tau(f) = \langle ev(f, 0), ev(f, 1), \ldots \rangle\).}

Starting with our \(\Pi\Sigma\)-field \((K(x), \sigma)\) over \(K\) and the evaluation function \(\{40\}\) we can construct for an AP-extension an appropriate evaluation function by iterative application of Lemma 2. In particular, this yields a \(K\)-homomorphism from the given AP-extension into the ring of sequences by Lemma 1.

Finally, we utilize the following result from [Sch17]; compare [OS18, Lemma 2.2].

Theorem 2. Let \((A, \sigma)\) be a difference field with constant field \(K\) and let \((E, \sigma)\) be a basic RI\(I\)-extension of \((A, \sigma)\). Then any \(K\)-homomorphism \(\tau : E \rightarrow S(K)\) is injective.

In other words, if we succeed in modeling our nested products within a basic RI\(I\)-extension (in particular, a multiple chain RI\(I\)-extension) over \((K(x), \sigma)\) with an appropriate evaluation function, then we automatically obtain a \(K\)-embedding.

### 4.2. A structural theorem for multiple chain \(\Pi\)-extensions

In part 1 of Theorem 1 a criterion is given that enables one to check, e.g., the algorithms from [Kar81, Sch16] whether a P-extension is a \(\Pi\)-extension. In [OS18, Lemma 5.1] (based on [Sch10a, Sch17]) this criterion has been generalized for “single nested” P-extensions as follows.

Lemma 5. Let \((E, \sigma)\) be a difference field and let \(f_1, \ldots, f_s \in E^*\). Then the following statements are equivalent.

1. There do not exist \((v_1, \ldots, v_s) \in \mathbb{Z}^s \setminus \{0_s\}\) and \(g \in E^*\) such that \(\frac{\sigma(g)}{g} = f_1^{v_1} \cdots f_s^{v_s}\) holds.
2. The P-extension \((E[z_1, z_1^{-1}] \ldots [z_s, z_s^{-1}], \sigma)\) of \((E, \sigma)\) with \(\sigma(z_i) = f_i z_i\) for \(1 \leq i \leq s\) is a \(\Pi\)-extension.
3. The difference field extension \((E(z_1) \ldots (z_s), \sigma)\) of \((E, \sigma)\) with \(z_i\) transcendental over \(E(z_1) \ldots (z_{i-1})\) and \(\sigma(z_i) = f_i z_i\) for \(1 \leq i \leq s\) is a \(\Pi\)-field extension.

In Theorem 3, we will extend this result further to multiple-chain P-extensions. Here we utilize that a solution for a certain class of homogeneous first-order difference equations has a particularly simple form; this result is a specialization of [Sch16, Corollary 4.6].

Corollary 1. Let \((E, \sigma)\) be a \(\Pi\)-extension of a difference field \((A, \sigma)\) with \(E = A \langle t_1 \rangle \ldots \langle t_c \rangle\). Then for any \(g \in E \setminus \{0\}\) with \(\sigma(g) = ut_1^{v_1} \cdots t_c^{v_c}\) for some \(u \in A^*\) and \(z_i \in \mathbb{Z}\) we have

\[ g = h t_1^{v_1} \cdots t_c^{v_c} \]

with \(h \in A^*\) and \(v_i \in \mathbb{Z}\).

Now we are ready to prove a general criterion that enables one to check if a multiple chain P-extension forms a \(\Pi\)-extension. This result will be heavily used within the next two sections.

Theorem 3. Let \((H, \sigma)\) be a difference field and let \((H_\ell, \sigma)\) with \(H_\ell = H_\langle t_1 \rangle \ldots \langle t_{s_\ell} \rangle\) for \(1 \leq \ell \leq m\) be single chain P-extensions of \((H, \sigma)\) over \(H\) with base \(c_\ell \in H^*\) where \(s_1 \geq s_2 \geq \cdots \geq s_m\). In particular, the automorphisms are given by

\[ \sigma(t_{\ell,k}) = \alpha_{\ell,k} t_{\ell,k} \quad \text{where} \quad \alpha_{\ell,k} = c_\ell t_{\ell,1} \cdots t_{\ell,k-1} \in (H^*)^{H_\langle t_{\ell,1} \rangle \ldots \langle t_{\ell,k-1} \rangle}. \]

Let \((A, \sigma)\) be the ordered multiple chain P-ext. of \((H, \sigma)\) with \(A = H_\langle t_{1,1} \rangle \ldots \langle t_{w_1,1} \rangle \ldots \langle t_{1,d} \rangle \ldots \langle t_{w_d,1} \rangle\) of monomial depth \(d := \max(s_1, \ldots, s_m)\) with \(m = w_1 \geq w_2 \geq \cdots \geq w_d\) composed by the single chain
Π-extensions \((\mathbb{H}, \sigma)\) of \((\mathbb{H}, \sigma)\) with the automorphism \([59]\). Then \((\mathbb{A}, \sigma)\) is a Π-extension of \((\mathbb{H}, \sigma)\) if and only if there does not exist a \(g \in \mathbb{H}^*\) and \((v_1, \ldots, v_m) \in \mathbb{Z}^m \setminus \{0_m\}\) such that
\[
\frac{\sigma(g)}{g} = c_1^{v_1} \cdots c_m^{v_m}.
\]

\textbf{Proof.} “ \(\Rightarrow\)” Suppose that \((\mathbb{A}, \sigma)\) is a Π-extension of \((\mathbb{H}, \sigma)\). Then, it is a tower of Π-extensions \((\mathbb{A}_i, \sigma)\) of \((\mathbb{H}, \sigma)\) where \(\mathbb{A}_i = \mathbb{A}_{i-1}(t_{1,i}, \ldots, t_{w_{i-1},i})\) for \(1 \leq i \leq d\) with \(\mathbb{A}_0 = \mathbb{H}\). Since \((\mathbb{A}_1, \sigma)\) is a Π-extension of \((\mathbb{H}, \sigma)\), it follows by Lemma 5 that there does not exist a \(g \in \mathbb{H}\) and \((v_1, \ldots, v_w) \in \mathbb{Z}^w \setminus \{0_w\}\) with \(w_1 = m\) such that \([60]\) holds.

“ \(\Leftarrow\)” Conversely, suppose that there does not exist a \(g \in \mathbb{H}\) and \((v_1, \ldots, v_w) \in \mathbb{Z}^w \setminus \{0_w\}\) with \(w_1 = m\) such that \([60]\) holds. Let \((\mathbb{A}_1, \sigma)\) with \(\mathbb{A}_1 = \mathbb{H}(t_{1,1}, \ldots, t_{w_{1,1}})\) be a Π-extension of \((\mathbb{H}, \sigma)\) with \(\sigma(t_{j,1}) = \alpha_{j,1} t_{j,1}\) for all \(1 \leq j \leq w_1\). By Lemma 5 \((\mathbb{A}_1, \sigma)\) is a Π-extension of \((\mathbb{H}, \sigma)\). Let \((\mathbb{A}_i, \sigma)\) with \(\mathbb{A}_i = \mathbb{A}_{i-1}(t_{1,i}, \ldots, t_{w_{i-1},i})\) be the multiple chain Π-extension of \((\mathbb{H}, \sigma)\) with \(\sigma(t_{1,i}) = \cdots = \sigma(t_{w_{i-1},i})\) for all \(1 \leq i \leq d\) with the automorphism \([59]\). Assume that \((\mathbb{A}_k, \sigma)\) is a Π-extension of \((\mathbb{H}, \sigma)\) for all \(1 \leq k \leq \delta\) with \(d > \delta \geq 1\) and that \((\mathbb{A}_{\delta+1}, \sigma)\) is not a Π-extension of \((\mathbb{A}_\delta, \sigma)\). Then by Lemma 5 we can take a \(g \in \mathbb{A}_\delta \setminus \{0\}\) and \((v_1, v_2, \ldots, v_{w_{\delta+1}}) \in \mathbb{Z}^{w_{\delta+1}} \setminus \{0_{w_{\delta+1}}\}\) such that
\[
\sigma(g) = \alpha_1^{v_1} \cdots \alpha_{w_{\delta+1}}^{v_{w_{\delta+1}}} g
\]
holds. By Corollary 1 it follows that \(g = h t_{1,1}^{v_{1,1}} \cdots t_{2,2}^{v_{2,2}} \cdots t_{w_{\delta+1},\delta}^{v_{w_{\delta+1},\delta}}\) with \(h \in \mathbb{H}^*\) and \(v_{i,j} \in \mathbb{Z}\). For the left hand side of \([61]\) we have that
\[
\sigma(g) = \gamma t_{1,\delta}^{v_{1,\delta}} \cdots t_{w_{\delta},\delta}^{v_{w_{\delta},\delta}}
\]
where \(\gamma \in \mathbb{A}_{\delta-1}\) and for the right hand side of \([61]\) we have that
\[
\alpha_1^{v_1} \cdots \alpha_{w_{\delta+1}}^{v_{w_{\delta+1}}} g = \omega t_{1,\delta}^{v_{1,\delta}} \cdots t_{w_{\delta+1},\delta}^{v_{w_{\delta+1},\delta}}
\]
where \(\omega \in \mathbb{A}_{\delta-1}\). Consequently, \(v_{k,\delta} = v_{k,\delta} + v_k\) and thus \(v_k = 0\) for all \(1 \leq k \leq w_{\delta+1}\) which is a contradiction to the assumption that \((v_1, \ldots, v_{w_{\delta+1}}) \neq 0_{w_{\delta+1}}\). Thus \((\mathbb{A}_d, \sigma)\) is a Π-extension of \((\mathbb{H}, \sigma)\).

\[\square\]

5. The main building blocks to represent nested products in ΠIΣ-extensions

Suppose that we are given a finite set of hypergeometric products of finite nesting depth which have been brought into the form as given in Proposition 2. In the following we will show in Sections 5.1 and 5.2 how these hypergeometric and geometric products can be modelled in ΠIΣ-extensions. For the treatment of geometric products one has to deal in addition with products defined over roots of unity of finite nesting depth. This extra complication will be treated in Section 5.3. Finally, in Section 6 below we will combine all these techniques to represent the full class of hypergeometric products of finite nesting depth in ΠIΣ-extensions.

5.1. Nested Hypergeometric products with shift-coprime multiplicands

In Proposition 2 we showed that a finite set of hypergeometric products of finite nesting depth can be brought in a shift-coprime product representation form. In the setting of ΠIΣ-field extensions the underlying Definition 4 can be generalized as follows.

\textbf{Definition 13.} Let \((\mathbb{F}(t), \sigma)\) be a ΠIΣ-field extension of \((\mathbb{F}, \sigma)\). We call two polynomials \(f, g \in \mathbb{F}[t]\) shift-coprime (or \(\sigma\)-coprime) if for all \(k \in \mathbb{Z}\) we have that \(\gcd(f, \sigma^k(h)) = 1\).

\footnote{Note that \((c_1, \ldots, c_m) = (\alpha_1, \ldots, \alpha_{w_{1,1}})\).}
Inspired by the key tool to represent hypergeometric products of nesting depth 1 in a \( \Pi \)-extension.

**Theorem 4.** Let \( (\mathbb{F}(t),\sigma) \) be a \( \Pi\Sigma \)-extension of \((\mathbb{F},\sigma)\). Let \( f_1,\ldots,f_s \in \mathbb{F}[t] \setminus \mathbb{F} \) be irreducible monic polynomials. Then the following statements are equivalent.

1. For all \( i, j \) with \( 1 \leq i < j \leq s \), \( f_i \) and \( f_j \) are shift-coprime.
2. There does not exist \( (v_1,\ldots,v_s) \in \mathbb{Z}^s \setminus \{0,1\} \) and \( g \in \mathbb{F}(t)^* \) with \( \frac{\sigma(g)}{g} = f_1^{v_1} \cdots f_s^{v_s} \).
3. The \( P \)-extension \( (\mathbb{F}(t)[z_1,z_1^{-1}]\cdots[z_s,z_s^{-1}],\sigma) \) of \((\mathbb{F}(t),\sigma)\) with \( \sigma(z_i) = f_i z_i \) for \( 1 \leq i \leq s \) is a \( \Pi \)-extension.

With this result we are now in the position to refine Theorem 3 in order to construct a \( \Pi \)-extension in which we can model hypergeometric products of finite nesting depth that are in shift-coprime representation form.

**Theorem 5.** Let \( (\mathbb{F}(t),\sigma) \) be a \( \Pi\Sigma \)-extension of \((\mathbb{F},\sigma)\). Let \( f = (f_1,\ldots,f_m) \in (\mathbb{F}[t] \setminus \mathbb{F})^m \) be irreducible monic polynomials. For all \( 1 \leq \ell \leq m \), let \( (\mathbb{F}_\ell,\sigma) \) with \( \mathbb{F}_\ell := \mathbb{F}(t)\langle z_{\ell,1}\rangle\cdots\langle z_{\ell,s}\rangle \) be a single chain \( P \)-extension of \((\mathbb{F}(t),\sigma)\) with base \( f_\ell \in \mathbb{F}[t] \setminus \mathbb{F} \) with the automorphism

\[
\sigma(z_{\ell,k}) = \alpha_{\ell,k} z_{\ell,k} \quad \text{where} \quad \alpha_{\ell,k} = f_\ell z_{\ell,1} \cdots z_{\ell,k-1} \in (\mathbb{F}^*)_{F}(z_{\ell,1})\cdots(z_{\ell,k-1})
\]

and with \( s_1 \geq s_2 \geq \cdots \geq s_m \). Let \((\mathbb{H}_b,\sigma)\) with

\[
\mathbb{H}_b = \mathbb{F}(t)\langle z_1\rangle\cdots\langle z_b \rangle = \mathbb{F}(t)\langle z_{1,1}\rangle\cdots\langle z_{1,s}\rangle\cdots\langle z_{b,1}\rangle\cdots\langle z_{b,s}\rangle
\]

be an ordered multiple chain \( P \)-extension of \((\mathbb{F}(t),\sigma)\) of monomial depth \( b = \max(s_1,\ldots,s_m) \) with bases \( f_1,\ldots,f_m \) where \( m = w_1 \geq w_2 \geq \cdots \geq w_b \) which is composed by the single chain \( P \)-extensions \((\mathbb{F}_\ell,\sigma)\) of \((\mathbb{F}(t),\sigma)\). Then \((\mathbb{H}_b,\sigma)\) is a \( \Pi \)-extension of \((\mathbb{F}(t),\sigma)\) if and only if for all \( i, j \) with \( 1 \leq i < j \leq m \) the \( f_i \) and \( f_j \) are shift-coprime.

**Proof.** "\( \implies \)" If \((\mathbb{H}_b,\sigma)\) is a \( \Pi \)-extension of \((\mathbb{F}(t),\sigma)\), then by Theorem 3 there does not exist \( g \in \mathbb{F}(t)^* \) such that \( \frac{\sigma(g)}{g} = f_1^{v_1} \cdots f_m^{v_m} \) holds, and by Theorem 4 for all \( i, j \) with \( 1 \leq i < j \leq m \), \( f_i \) and \( f_j \) are shift-coprime.

"\( \impliedby \)" Conversely, if for all \( i, j \) with \( 1 \leq i < j \leq m \), \( f_i \) and \( f_j \) are shift-coprime, then by Theorem 4 there does not exist \( g \in \mathbb{F}(t)^* \) such that \( \frac{\sigma(g)}{g} = f_1^{v_1} \cdots f_m^{v_m} \) holds, and by Theorem 3 \((\mathbb{H}_b,\sigma)\) is a \( \Pi \)-extension of \((\mathbb{F}(t),\sigma)\).

Summarizing, we obtain the following crucial result.

**Corollary 2.** Let \((\mathbb{K}(x),\sigma)\) be a rational difference field with the automorphism \( \sigma(x) = x + 1 \) and the evaluation function \( \text{ev} : \mathbb{K}(x) \times \mathbb{Z}_{\geq 0} \rightarrow \mathbb{K} \) given by \( \{40\} \). Let \( H_1(n),\ldots,H_e(n) \) be hypergeometric products in \( \text{Prod}_{\text{d}}(\mathbb{K}(x)) \) of nesting depth at most \( b \) which are all in shift-coprime representation form (see Definition 3) and which are all \( \delta \)-refined for some \( \delta \in \mathbb{Z}_{\geq 0} \). Then one can construct an ordered multiple chain \( \Pi \)-extension \((\mathbb{H}_\delta,\sigma)\) of \((\mathbb{K}(x),\sigma)\) with

\[
\mathbb{H}_\delta = \mathbb{K}(x)\langle z_1\rangle\cdots\langle z_b \rangle = \mathbb{K}(x)\langle z_{1,1}\rangle\cdots\langle z_{1,s}\rangle\cdots\langle z_{b,1}\rangle\cdots\langle z_{b,s}\rangle
\]

which is composed by the single chain \( \Pi \)-extensions \((\mathbb{F}_\ell,\sigma)\) of \((\mathbb{K}(x),\sigma)\) with \( \mathbb{F}_\ell = \mathbb{K}(x)\langle z_{\ell,1}\rangle\cdots\langle z_{\ell,s}\rangle \) with

1. the automorphism \( \sigma : \mathbb{F}_\ell \rightarrow \mathbb{F}_\ell \) defined by

\[
\sigma(z_{\ell,k}) = \tilde{\alpha}_{\ell,k} z_{\ell,k} \quad \text{where} \quad \tilde{\alpha}_{\ell,k} = f_\ell z_{\ell,1} \cdots z_{\ell,k-1} \in (\mathbb{K}(x)^*)_{F}(z_{\ell,1})\cdots(z_{\ell,k-1})
\]

for \( 1 \leq \ell \leq p_1 \) and \( 1 \leq k \leq s_\ell \) where \( f_\ell \in \mathbb{K}[x] \setminus \mathbb{K} \) is an irreducible monic polynomial, and
(2) the evaluation function $\tilde{ev}: \tilde{F}_\ell \times \mathbb{Z}_{\geq 0} \to \mathbb{K}$ given by $\tilde{ev}|_{\mathbb{K}(x)} = ev$ with (66) and

$$
\tilde{ev}(z_{\ell,k}, n) = \prod_{j=\delta}^{n} \tilde{ev}(\tilde{a}_{\ell,k}, j-1)
$$

for $1 \leq \ell \leq p_1$ and $1 \leq k \leq s_\ell$

with the following property: for all $1 \leq i \leq e$ there are $k, j$ such that

$$
ev(z_{k,j}, n) = \tilde{H}_i(n), \quad \forall n \geq \max(0, \delta - 1).
$$

Furthermore, for all $g \in \mathbb{H}_b$, the map $\tilde{\tau}: \mathbb{H}_b \to S(\mathbb{K})$ defined by

$$\tilde{\tau}(g) = (\tilde{ev}(g, n))_{n \geq 0}
$$

is a $\mathbb{K}$-embedding. If $\mathbb{K}$ is computable, the above construction can be given explicitly.

**Proof.** By the procedure outlined in Remark 4 (skipping step (1) since the input is already in the right form) we obtain the ordered multiple chain P-extension $(\tilde{\mathbb{H}}_b, \sigma)$ of $(\mathbb{K}(x), \sigma)$ with (63) and (65) such that (66) holds for all $1 \leq i \leq e$ for some $j, k$. Since the bases $\tilde{f}_1, \ldots, \tilde{f}_{p_1}$ of the single chain P-extensions $(\tilde{\mathbb{F}}_1, \sigma), \ldots, (\tilde{\mathbb{F}}_{p_1}, \sigma)$ that composes $(\tilde{\mathbb{H}}_b, \sigma)$ are shift-coprime, it follows by Theorem 5 that $(\tilde{\mathbb{H}}_b, \sigma)$ is a II-extension of $(\mathbb{K}(x), \sigma)$. Since $(\tilde{\mathbb{H}}_b, \sigma)$ is a basic II-extension of the rational difference field $(\mathbb{K}(x), \sigma)$, it follows by Theorem 2 that the $\mathbb{K}$-homomorphism $\tilde{\tau}: \mathbb{H}_b \to S(\mathbb{K})$ defined by (67) is injective. Since $\mathbb{K}$ is computable, all the above ingredients can be constructed explicitly. \hfill \Box

**Example 8 (Cont. Example 7).** Consider the ordered multiple chain P-extension $(\tilde{\mathbb{H}}, \sigma)$ of the rational difference field $(\mathbb{K}(x), \sigma)$ of monomial depth 2 with $\tilde{\mathbb{H}} = \mathbb{K}(x)(z_{1,1})(z_{2,1})(z_{1,2})$ where $(\mathbb{H}, \sigma)$ is composed by the single chain II-extensions of $(\mathbb{K}(x), \sigma)$ constructed in parts (7) and (8) of Example 7. Since the bases of $(\tilde{\mathbb{H}}, \sigma)$ given by $(x - 2)$ and $(x + \frac{1}{24})$ are shift-coprime with respect to the automorphism $\sigma(x) = x + 1$, it follows that the ordered multiple chain P-extension $(\tilde{\mathbb{H}}, \sigma)$ of the rational difference field $(\mathbb{K}(x), \sigma)$ of monomial depth 2 is a II-extension. Furthermore, it follows by Theorem 2 that the map $\tilde{\tau}: \mathbb{H} \to S(\mathbb{K})$ defined by $\tilde{\tau}(f) = (\tilde{ev}(f, n))_{n \geq 0}$ for all $f \in \mathbb{H}$ is a $\mathbb{K}$-embedding where $\tilde{ev} = ev$ defined in (54) and (55). In particular, for the expression $g$ given by (57) we have that $\tilde{H}(n) = \tilde{ev}(g, n)$ holds for all $n \geq 2$.

5.2. Geometric products

In Karr’s algorithm [Kar81] and all the improvements [KS06, Sch07a, Sch08, AP10, Sch15, Sch16, Sch17] one relies on certain algorithmic properties of the constant field $\mathbb{K}$. Among those, one needs to solve the following problem.

**Problem GO for $\alpha_1, \ldots, \alpha_w \in K^*$**

Given a field $K$ and $\alpha_1, \ldots, \alpha_w \in K^*$. Compute a basis of the submodule

$$V := \{(u_1, \ldots, u_w) \in \mathbb{Z}^w \mid \prod_{i=1}^{w} \alpha_i^{u_i} = 1\} \text{ of } \mathbb{Z}^w \text{ over } \mathbb{Z}.$$

In our approach Problem GO is crucial to solve Problem RPE, but one has to solve it not only in a given field $K$ (compare the definition of $\sigma$-computable in [Sch05, KS06]) but one must be able to solve it in any finite algebraic field extension of $K$. This gives rise to the following definition.

**Definition 14.** A field $K$ is strongly $\sigma$-computable if the standard operations in $K$ can be performed, multivariate polynomials can be factored over $K$ and Problem GO can be solved for $K$ and any finite algebraic field extension of $K$.  

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Note that Ge’s algorithm [Ge93a] (see also [Kau06, Algorithm 7.16, page 84]) solves Problem G0 over an arbitrary number field $K$. Since any finite algebraic extension of an algebraic number field is again an algebraic number field, we obtain the following result; for a weaker result see [Sch05, Theorem 3.5].

**Lemma 6.** An algebraic number field $K$ is strongly $\sigma$-computable.

By [OS18, Theorem 5.4] and the consideration of [OS18, pg. 204] (see also [Oca19, Lemma 5.2.2]) we provided an algorithm that enabled us to handle geometric products of nested depth 1. More precisely, given a $P$-extension that models such products, Lemma 7 states that one can construct an RII-extension in which the products can be rephrased.

**Lemma 7.** Let $\mathbb{K} = K(\kappa_1, \ldots, \kappa_n)$ be a rational function field over a field $K$ and $(\mathbb{K}, \sigma)$ be a difference field with $\sigma(c) = c$ for all $c \in \mathbb{K}$. Let $\langle K(x_1) \ldots \rangle_{x_e}, \sigma)$ be a $P$-extension of $(\mathbb{K}, \sigma)$ with $\sigma(x_i) = \gamma_i x_i$ for $1 \leq i \leq e$ where $\gamma_i \in \mathbb{K}^*$. Let $ev : K(x_1) \ldots \langle x_e \rangle \times \mathbb{Z}_{\geq 0} \rightarrow \mathbb{K}$ be the evaluation function defined by $ev(x_i, n) = \gamma_i^n$ for $1 \leq i \leq e$. Then:

1. One can construct an RII-extension $\langle \mathbb{K}(\vartheta) \langle \tilde{y}_1 \rangle \ldots \langle \tilde{y}_s \rangle, \sigma \rangle$ of $(\mathbb{K}, \sigma)$ with

   $\sigma(\vartheta) = \zeta \vartheta$ and $\sigma(\tilde{y}_k) = \alpha_k \tilde{y}_k$ \hspace{1cm} (68)

   for $1 \leq k \leq s$ where $\bar{K} = \bar{K}(\kappa_1, \ldots, \kappa_n)$ and $\bar{K}$ is a finite algebraic field extension of $K$ with $\zeta \in \bar{K}$ being a primitive $\lambda$-th root of unity and $\alpha_k \in \bar{K}^*$;

2. one can construct the evaluation function $\tilde{ev} : \mathbb{K}(\vartheta) \langle \tilde{y}_1 \rangle \ldots \langle \tilde{y}_s \rangle \times \mathbb{Z}_{\geq 0} \rightarrow \bar{K}$ defined as

   $\tilde{ev}(\vartheta, n) = \zeta^n$ and $\tilde{ev}(\tilde{y}_k, n) = \alpha_k^n$; \hspace{1cm} (69)

3. one can construct a difference ring homomorphism $\varphi : K(x_1) \ldots \langle x_e \rangle \rightarrow \bar{K}(\vartheta) \langle \tilde{y}_1 \rangle \ldots \langle \tilde{y}_s \rangle$ with

   $\varphi(x_i) = \vartheta^{\mu_i} \tilde{y}_i^{\nu_i} = \vartheta^{\mu_i} \tilde{y}_i^{\nu_i - 1} \ldots \tilde{y}_s^{\nu_i - s}$ \hspace{1cm} (70)

   for $1 \leq i \leq e$ where $0 \leq \mu_i < \lambda$ and $v_i, k \in \mathbb{Z}$ for $1 \leq k \leq s$

   such that for all $f \in K(x_1) \ldots \langle x_e \rangle$ and for all $n \in \mathbb{Z}_{\geq 0}$,

   $ev(f, n) = \tilde{ev}(\varphi(f), n)$

holds. If $K$ is strongly $\sigma$-computable, then the above constructions are computable.

Using this result we will derive an extended version in Lemma 9 that deals with the class of ordered multiple chain AP-extensions that models geometric products of arbitrary but finite nesting depth.

In the following let $m \in \mathbb{Z}_{\geq 1}$, and for $1 \leq \ell \leq m$ let $(K_{\ell}, \sigma)$ with $K_{\ell} = K(y_{\ell}) = K(y_{\ell,1}, \ldots, y_{\ell,s_{\ell}})$ be a single chain P-extension of $(K, \sigma)$ with base $h_{\ell} \in \mathbb{K}$ where

$\sigma(y_{\ell,i}) = \alpha_{\ell,i} y_{\ell,i}$ with $\alpha_{\ell,i} = h_{\ell} y_{\ell,1} \ldots y_{\ell,i-1} \in (K^*)_{K}(y_{\ell,1}, \ldots, y_{\ell,i-1})$. \hspace{1cm} (71)

In particular, we assume that $s_1 \geq s_2 \geq \cdots \geq s_m$. Let $ev : K_{\ell} \times \mathbb{Z}_{\geq 0} \rightarrow \mathbb{K}$ be the evaluation function defined by

$ev(y_{\ell,i}, n) = \prod_{j=1}^{n} ev(\alpha_{\ell,i}, j-1) = \prod_{j=1}^{n} \alpha_{\ell,i}$; \hspace{1cm} (72)

---

"For concrete instances the R-monomial $\vartheta$ might be obsolete. In particular, if $\mu_i = 0$ for all $1 \leq i \leq e$ in (70) it can be removed."
in particular, for all \( c \in \mathbb{K} \) and \( n \geq 0 \) we set \( \text{ev}(c, n) = c \). Let \((\mathbb{A}, \sigma)\) be the multiple chain P-extension of \((\mathbb{K}, \sigma)\) built by the single chain\( \Pi \)-extensions \((\mathbb{K}_\ell, \sigma)\) of \((\mathbb{K}, \sigma)\) over \(\mathbb{K}\). That is,

\[
\mathbb{A} = \mathbb{K} \langle y_1 \rangle \langle y_2 \rangle \cdots \langle y_m \rangle = \mathbb{K} \langle y_{1,1} \rangle \cdots \langle y_{1,s_1} \rangle \langle y_{2,1} \rangle \cdots \langle y_{2,s_2} \rangle \cdots \langle y_{m,1} \rangle \cdots \langle y_{m,s_m} \rangle.
\]

We emphasize that all the \( y_{\ell,i} \) model 1-refined geometric products in product factored form of an arbitrary but finite nesting depth. Depending on the context, \( y_{\ell} \) denotes \( (y_{1,1}, \ldots, y_{1,s_1}) \) or \( y_{1,1}, \ldots, y_{\ell,s_\ell} \) or \( y_{1,1} \cdots y_{\ell,s_\ell} \). Note that the P-monomials \( y_{\ell,i} \) can be ordered in increasing order of their depths. Namely, take the depth function \( \delta : \mathbb{A} \to \mathbb{Z}_{\geq 0} \) over \(\mathbb{K}\) of \((\mathbb{A}, \sigma)\) and let \( d = \max(s_1, s_2, \ldots, s_m) \) be the maximal depth. Then taking \( \mathbb{A}_0 = \mathbb{K} \) we can consider the tower of P-extensions \((\mathbb{A}_i, \sigma)\) of \((\mathbb{A}_{i-1}, \sigma)\) with

\[
A_i = \mathbb{A}_{i-1} \langle y_{1,i} \rangle = \mathbb{A}_{i-1} \langle y_{2,i} \rangle \cdots \langle y_{w_i,i} \rangle
\]

for \( 1 \leq i \leq d \) where \( m = w_1 \geq w_2 \geq \cdots \geq w_d \) and with the automorphism \((71)\) for \( 1 \leq \ell \leq w_i \). In this way, the P-monomials at the \( i \)-th extension have the depth \( \delta(y_{1,i}) = \delta(y_{2,i}) = \cdots = \delta(y_{w_i,i}) = i \). Further, the ring \( \mathbb{A}_d \) is isomorphic to \( \mathbb{A} \) up to reordering of the P-monomials. In particular, \((\mathbb{A}_d, \sigma_d)\) is an ordered multiple chain P-extension of \((\mathbb{K}, \sigma)\) of monomial depth at most \( d \) induced by the single chain\( \Pi \)-extensions \((\mathbb{K}_\ell, \sigma)\) of \((\mathbb{K}, \sigma)\) for \( 1 \leq \ell \leq m \) with \((71)\) and \((72)\). Observe that since \( \mathbb{A}_d \simeq \mathbb{A} \), the evaluation function \( \text{ev} : \mathbb{A}_i \times \mathbb{Z}_{\geq 0} \to \mathbb{K} \) for all \( i \) with \( 1 \leq i \leq d \) is also defined by \((72)\).

In order to derive the main result of this subsection in Lemma\( 9 \) we need following simple construction.

**Lemma 8.** Let \((\mathbb{A}(t), \sigma)\) be a \( \Pi \)-extension of \((\mathbb{A}, \sigma)\) with \( \sigma(t) = \alpha t \) and let \((\mathbb{H}, \sigma)\) be a difference ring. Let \( \rho : \mathbb{A} \to \mathbb{H} \) be a difference ring homomorphism and let \( \rho : \mathbb{A}(t) \to \mathbb{H} \) be a ring homomorphism defined by \( \rho|_{\mathbb{A}} = \tilde{\rho} \) and \( \rho(t) = g \) for some \( g \in \mathbb{H} \). If \( \sigma(g) = \rho(\alpha)g \), then \( \rho \) is a difference ring homomorphism.

**Proof.** Suppose that \( \sigma(g) = \rho(\alpha)g \) holds. Then \( \sigma(\rho(f)) = \sigma(g) = \rho(\alpha)g = \rho(\alpha t) = \rho(\sigma(t)) \). Consequently, \( \sigma(\rho(f)) = \rho(\sigma(f)) \) for all \( f \in \mathbb{A}(t) \). \( \square \)

**Lemma 9.** For \( 1 \leq \ell \leq m \), let \((\mathbb{K}_\ell, \sigma)\) with \( \mathbb{K}_\ell = \mathbb{K}(y_{1,1}, \ldots, y_{\ell,s_\ell}) \) be single chain P-extensions of \((\mathbb{K}, \sigma)\) over a rational function field \( \mathbb{K} = K(\kappa_1, \ldots, \kappa_u) \) with base \( h_\ell \in \mathbb{K}^* \), the automorphisms \((71)\) and the evaluation functions \((72)\).

Let \( d := \max(s_1, \ldots, s_m) \) and \( \mathbb{A}_0 = \mathbb{K} \). Consider the tower of difference ring extensions \((\mathbb{A}_i, \sigma)\) of \((\mathbb{A}_{i-1}, \sigma)\) where \( \mathbb{A}_i = \mathbb{A}_{i-1} \langle y_{1,i} \rangle \langle y_{2,i} \rangle \cdots \langle y_{w_i,i} \rangle \) for \( 1 \leq i \leq d \) with \( m = w_1 \geq \cdots \geq w_d \), the automorphism \((72)\) and the evaluation function \((72)\). In particular, one gets \((\mathbb{A}_d, \alpha)\) as the ordered multiple chain P-extension of \((\mathbb{K}, \sigma)\) of monomial depth at most \( d \) composed by the single chain P-extensions \((\mathbb{K}_\ell, \sigma)\) of \((\mathbb{K}, \sigma)\) for \( 1 \leq \ell \leq m \) with \((71)\) and \((72)\). Then one can construct

(a) an ordered multiple chain \( \Lambda \)-extension \((\mathbb{G}_d, \sigma)\) of \((\mathbb{K}, \sigma)\) of monomial depth at most \( d \) with \( \mathbb{K} = \mathbb{K}(\kappa_1, \ldots, \kappa_u) \) where \( \mathbb{K} \) is a finite algebraic extension of \( K \), with

\[
\mathbb{G}_d = \mathbb{K}(\vartheta_{1,1}) \cdots \langle \vartheta_{1,d} \rangle \langle \tilde{y}_{1,1} \rangle \cdots \langle \tilde{y}_{1,d} \rangle \langle \vartheta_{2,1} \rangle \cdots \langle \vartheta_{2,d} \rangle \langle \tilde{y}_{2,1} \rangle \cdots \langle \tilde{y}_{2,d} \rangle \langle \vartheta_{3,1} \rangle \cdots \langle \vartheta_{3,d} \rangle \cdots \langle \vartheta_{w_d,1} \rangle \cdots \langle \vartheta_{w_d,d} \rangle \langle \tilde{y}_{w_d,1} \rangle \cdots \langle \tilde{y}_{w_d,d} \rangle
\]

\[ (73) \]

where [3] \( v \geq 0, v_i \geq 0 \). Here the automorphism is defined for the \( \Lambda \)-monomials by

\[
\sigma(\vartheta_{1,k}) = \gamma_{1,k} \vartheta_{1,k} \quad \text{where} \quad \gamma_{1,k} = \zeta^{d_{1,k}} \vartheta_{1,1} \cdots \vartheta_{1,k-1} \in \mathbb{K}_\mathbb{G}(\vartheta_{1,1} \cdots \vartheta_{1,k-1})
\]

\[ (74) \]

for \( 1 \leq k \leq d \) and \( 1 \leq \ell \leq \ell_k \) where \( \mathbb{U} = \mathbb{U}(\zeta) \) is a multiplicative cyclic subgroup of \( \mathbb{K}^* \), generated by a primitive \( \lambda \)-th root of unity \( \zeta \in \mathbb{K}^* \), and the automorphism is defined for the \( \Lambda \)-monomials by

\[
\sigma(\tilde{y}_{\ell,k}) = \tilde{\alpha}_{\ell,k} \tilde{y}_{\ell,k} \quad \text{where} \quad \tilde{\alpha}_{\ell,k} = \tilde{h}_{\ell,k} \tilde{y}_{\ell,1} \cdots \tilde{y}_{\ell,k-1} \in \mathbb{K}(\tilde{h}_{\ell,1} \cdots \tilde{h}_{\ell,k-1})
\]

\[ (75) \]

for \( 1 \leq k \leq d \) and \( 1 \leq \ell \leq \ell_k \); [Note that if \( v_i = 0 \) or \( e_i = 0 \), then there is no depth-\( i \) \( \Lambda \)-monomial or \( \Lambda \)-monomial of depth \( i \), respectively.]

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(b) an evaluation function \( \hat{\mathrm{ev}} : G_d \times \mathbb{Z}_{\geq 0} \to \hat{\mathbb{K}} \) defined by

\[
\hat{\mathrm{ev}}(\hat{y}_{\ell,k}, n) = \prod_{j=1}^{n} \mathrm{ev}(\gamma_{\ell,k}, j - 1) \quad \text{and} \quad \hat{\mathrm{ev}}(\hat{y}_{\ell,k}, n) = \prod_{j=1}^{n} \hat{\mathrm{ev}}(\hat{\alpha}_{\ell,k}, j - 1);
\]

(76)

(c) a difference ring homomorphism \( \rho_d : \hat{A}_d \to G_d \) defined by \( \rho_d|_{\mathbb{K}} = \text{id}_{\mathbb{K}} \) and

\[
\rho_d(\hat{y}_{\ell,k}) = \hat{y}_{\ell,k}^{\mu_{\ell,k}} \hat{y}_{\ell,k}^{\nu_{\ell,k}} = \hat{y}_{\ell,k}^{\nu_{\ell,k}} \hat{y}_{\ell,k}^{\nu_{\ell,k}} \hat{y}_{\ell,k}^{\nu_{\ell,k}} \cdots \hat{y}_{\ell,k}^{\nu_{\ell,k}}
\]

for \( 1 \leq \ell \leq m \) and \( 1 \leq k \leq s_{\ell} \) with \( \mu_{\ell,i,k} \in \mathbb{Z}_{\geq 0} \) for \( 1 \leq i \leq v_k \) and \( v_{\ell,i,k} \in \mathbb{Z} \) for \( 1 \leq i \leq e_k \) such that the following properties hold:

1. There does not exist a \( (v_1, \ldots, v_{e_1}) \in \mathbb{Z}^{e_1} \setminus \{0_{e_1}\} \) with \( h_{e_1}^{v_1} \cdots h_{e_1}^{v_{e_1}} = 1 \).

2. The \( \mathbb{P} \)-extension \( (\hat{A}_d, \sigma) \) of \( (\mathbb{K}, \sigma) \) with

\[
\hat{A}_d = \mathbb{K} \langle \hat{y}_1 \rangle \langle \hat{y}_2 \rangle \cdots \langle \hat{y}_d \rangle = \mathbb{K} \langle \hat{y}_{1,1} \rangle \langle \hat{y}_{2,1} \rangle \cdots \langle \hat{y}_{e_1,1} \rangle \langle \hat{y}_{1,2} \rangle \cdots \langle \hat{y}_{e_2,2} \rangle \cdots \langle \hat{y}_{e_d,d} \rangle
\]

and the automorphism given in (75) is a \( \Pi \)-extension. In particular, it is an ordered multiple chain \( \Pi \)-extension of monomial depth \( d \).

3. For all \( f \in \hat{A}_d \) and for all \( n \in \mathbb{Z}_{\geq 0} \) we have

\[
\hat{\mathrm{ev}}(f, n) = \hat{\mathrm{ev}}(\rho_d(f), n).
\]

If \( K \) is strongly \( \sigma \)-computable, then the above constructions are computable.

**Proof.** Let \( (\hat{A}_d, \sigma) \) with \( \hat{A}_d = \hat{A}_{d-1} \langle \hat{y}_{1,d} \rangle \langle \hat{y}_{2,d} \rangle \cdots \langle \hat{y}_{w_d,d} \rangle \) be the ordered multiple chain \( \mathbb{P} \)-extension of \( (\mathbb{K}, \sigma) \) of monomial depth \( d \in \mathbb{Z}_{\geq 0} \) as described above with the automorphism (72) and the evaluation function (76) and the monomial depth \( d \).

If \( d = 1 \), statements (2) and (3) of the Lemma hold by Lemma 7. Hence by Lemma 7, there are no \( g \in \hat{\mathbb{K}}^* \) and \( (v_1, \ldots, v_{e_1}) \in \mathbb{Z}^{e_1} \setminus \{0_{e_1}\} \) with \( h_{e_1}^{v_1} \cdots h_{e_1}^{v_{e_1}} = \frac{\sigma(g)}{g} = 1 \) and thus also statement (1) of the Lemma holds.

Now let \( d \geq 2 \) and suppose that the Lemma holds for \( d - 1 \). That is, we can construct \( (G_{d-1}, \sigma) \) with

\[
G_{d-1} = \mathbb{K} \langle \hat{y}_{1,1} \rangle \cdots \langle \hat{y}_{e_2,1} \langle \hat{y}_{1,2} \rangle \cdots \langle \hat{y}_{e_2,2} \rangle \cdots \langle \hat{y}_{e_{d-1},d-1} \rangle \langle \hat{y}_{e_{d-1},d} \rangle
\]

which is an ordered multiple chain \( \mathbb{P} \)-extension of \( (\mathbb{K}, \sigma) \) of monomial depth \( d - 1 \) with the automorphism given by (74) for \( 1 \leq k \leq d - 1 \) and \( 1 \leq \ell \leq v_k \) and given by (75) for \( 1 \leq k \leq d - 1 \) and \( 1 \leq \ell \leq e_k \). In addition, we get the evaluation function \( \hat{\mathrm{ev}} : G_{d-1} \times \mathbb{Z}_{\geq 0} \to \hat{\mathbb{K}} \) defined as (76) and the difference ring homomorphism \( \rho_{d-1} : \hat{A}_{d-1} \to G_{d-1} \) defined by \( \rho_{d-1}|_{\mathbb{K}} = \text{id}_{\mathbb{K}} \) and (77) such that statements (1) and (2) of the Lemma hold. We prove the Lemma for the ordered multiple chain \( \mathbb{P} \)-extension \( (A_{d-1}, \sigma) \) of \( (\mathbb{K}, \sigma) \) with \( A_{d-1} = A_{d-1} \langle y_{1,d-1} \rangle \langle y_{2,d-1} \rangle \cdots \langle y_{w_d-1,d} \rangle \) where \( \delta(y_{1,d}) = \cdots = \delta(y_{w_d,d}) = d \).

Since the shift quotient of these \( \mathbb{P} \)-monomials is contained in \( A_{d-1}^* \), i.e.,

\[
\frac{\sigma(y_{1,d})}{y_{1,d}} = \alpha_{t,d} \in A_{d-1}^*,
\]

we can iteratively apply the difference ring homomorphism \( \rho_{d-1} : A_{d-1} \to G_{d-1} \) to rephrase each \( \alpha_{t,d} \) in \( G_{d-1} \). In particular, by Remark 3 we have \( \sigma^{-1}(\alpha_{t,d}) = y_{t,d-1} \) and thus by (77) we get

\[
h_{t,d} := \rho_{d-1}(\sigma^{-1}(\alpha_{t,d})) = \rho_{d-1}(y_{t,d-1}) = \rho_{d-1}(y_{t,d-1}) = \rho_{d-1}^{\mu_{d-1}} \rho_{d-1}^{\nu_{d-1}}
\]

(80)

\[\text{For all } c \in \mathbb{K}, \text{ we set } \hat{\mathrm{ev}}(c,n) = c \text{ for all } n \geq 0.\]

\[\text{Note that any } \mathbb{P} \text{-monomial } y_{\ell,k} \text{ with depth } k \text{ is mapped to a power product of } \mathbb{P} \text{-monomials having all depth } k.\]
where $\vartheta_{1,d-1}^{\mu_1,d-1} = \vartheta_{1,d-1}^{\mu_1,d-1} \cdots \vartheta_{v_{d-1},d-1}^{\mu_1,d-1}$ and $\vartheta_{d-1}^{\mu_1,d-1} = \vartheta_{1,d-1}^{\mu_1,d-1} \cdots \vartheta_{v_{d-1},d-1}^{\mu_1,d-1}$ for $1 \leq \ell \leq w_d$ with $\mu_{\ell,k,d-1} \in \mathbb{Z}_{\geq 0}$ for $1 \leq k \leq v_{d-1}$ and $v_{k,d-1} \in \mathbb{Z}$ for $1 \leq k \leq e_{d-1}$.

If $h_{\ell,d} = 1$, it follows with (79) ($d$ replaced by $d - 1$) that for all $n \in \mathbb{Z}_{\geq 0}$ we have

$$\text{ev}(y_{\ell,d}, n) = \prod_{j=1}^n \text{ev}(\alpha_{\ell,d}, j - 1) = \prod_{j=1}^n \text{ev}(\sigma^{-1}(\alpha_{\ell,d}), j) = \prod_{j=1}^n \sigma \text{ev}(\rho(\sigma^{-1}(\alpha_{\ell,d})), j) = \prod_{j=1}^n \sigma \text{ev}(h_{\ell,d}, j) = 1.$$

In particular, if $h_{\ell,d} = 1$ holds for all $1 \leq \ell \leq w_d$, we can set $G_{d} := G_{d-1}$ and extend $\rho_d$ to $\rho_d : A_d \rightarrow G_{d-1}$ with $\rho_d(y_{\ell,d}) = 1$ for $1 \leq \ell \leq w_d$. Thus the lemma hold.

Otherwise, take all AP-monomials in (80) for $1 \leq \ell \leq w_d$ with non-zero integer exponents. Then they belong to at least one of the single chain AP-extensions of $(\bar{K}, \sigma)$ in $(G_{d-1}, \sigma)$. Suppose there are $e_d \geq 0$ of these single chains II-extensions and $v_d \geq 0$ of them that are single chain A-extensions $(H_{b}, \sigma)$; note that we have $e_d + v_d \geq 1$. By appropriate reordering of $(G_{d-1}, \sigma)$ we may suppose that these $e_d$ single chain II-extensions $(F_{r}, \sigma)$ of $(\bar{K}, \sigma)$ with $1 \leq r \leq e_d$ are given by $F_r = \bar{K} \langle \bar{y}_{r,1} \rangle \langle \bar{y}_{r,2} \rangle \cdots \langle \bar{y}_{r,d-1} \rangle \langle \bar{y}_{r,d} \rangle$ and the $v_d$ A-extensions $(H_{b}, \sigma)$ of $(\bar{K}, \sigma)$ with $1 \leq b \leq v_d$ can be given by $H_b = \bar{K} \langle \partial_{b,1} \rangle \langle \partial_{b,2} \rangle \cdots \langle \partial_{b,d-1} \rangle \langle \partial_{b,d} \rangle$. Hence adjoin the P-monomials $\tilde{y}_{r,d}$ to $F_r$ with (74) where $k = d$ and $\ell = r$ yielding the single chain P-extensions $(F'_r, \sigma)$ of $(\bar{K}, \sigma)$ of monomial depth $d$ where

$$F'_r = F_r(\tilde{y}_{r,d}) = \bar{K} \langle \tilde{y}_{r,1} \rangle \langle \tilde{y}_{r,2} \rangle \cdots \langle \tilde{y}_{r,d-1} \rangle \langle \tilde{y}_{r,d} \rangle$$

and adjoin the A-monomial $\partial_{b,d}$ with (75) where $k = d$ and $\ell = b$ yielding the single chain A-extensions $(H'_b, \sigma)$ of $(\bar{K}, \sigma)$ of monomial depth $d$ where

$$H'_b = H_b(\partial_{b,d}) = \bar{K} \langle \partial_{b,1} \rangle \langle \partial_{b,2} \rangle \cdots \langle \partial_{b,d-1} \rangle \langle \partial_{b,d} \rangle.$$

Furthermore extend the evaluation functions $\sigma : F'_d \times \mathbb{Z}_{\geq 0} \rightarrow \bar{K}$ and $\sigma_v : H'_b \times \mathbb{Z}_{\geq 0} \rightarrow \bar{K}$ with (76) where $k = d$, $\ell = r$ or $k = d$, $\ell = b$, respectively. Now consider the multiple chain P-extension $(\bar{K}_d, \sigma)$ with

$$\bar{K}_d = \bar{K} \langle \bar{y}_{1,1} \rangle \cdots \langle \bar{y}_{1,d-1} \rangle \langle \bar{y}_{e_d-1,d-1} \rangle \langle \bar{y}_{e_d,d} \rangle \langle \tilde{y}_{d,d} \rangle$$

which one gets by taking all P-monomials in $G_{d-1}$ and the new P-monomials in $F'_r = F_r(\tilde{y}_{r,d})$ with $1 \leq r \leq e_d$. Here the automorphism is given by (74) and (75) and the equipped evaluation function is given by (76). Since there does not exist a $g \in \bar{K}$ and $(v_1, \ldots, v_d) \in \mathbb{Z}^d \setminus \{0_d\}$ with $\frac{\sigma(g)}{g} = h_1^{v_1} \cdots h_d^{v_d}$, it follows by Theorem 3 that $(\bar{K}_d, \sigma)$ is a II-extension of $(\bar{K}, \sigma)$. In particular, it is an ordered multiple chain II-extension of monomial depth $d$ by construction. Thus statements (1) and (2) of the Lemma hold. Finally, take the AP-extension $(G_d, \sigma)$ of $(G_{d-1}, \sigma)$ with

$$G_d = G_{d-1} \langle \tilde{y}_{1,1,d} \rangle \cdots \langle \tilde{y}_{d,d} \rangle \langle \tilde{y}_{e_d,d} \rangle \cdots \langle \tilde{y}_{e_d,d} \rangle = \bar{K} \langle \tilde{y}_{1,1,d} \rangle \cdots \langle \tilde{y}_{d,d} \rangle \langle \tilde{y}_{e_d,d} \rangle \cdots \langle \tilde{y}_{e_d,d} \rangle.$$

Let $1 \leq \ell \leq w_d$ and consider $h_{\ell,d}$ in (80). If $h_{\ell,d} = 1$, we define $g_{\ell} = 1$. In this case, $1 = \rho(h_{\ell,d}) = \rho_{d-1}(\sigma^{-1}(\alpha_{\ell,d})) = \sigma^{-1}(\rho_{d-1}(\alpha_{\ell,d}))$, and thus $\sigma(\frac{g_{\ell}}{g_{\ell}}) = 1 = \rho_{d-1}(\alpha_{\ell,d})$ holds. Otherwise, if $h_{\ell,d} \neq 1$, we set

$$g_{\ell} := \vartheta_{1,d-1}^{\mu_1,d-1} \cdots \vartheta_{v_{d-1},d-1}^{\mu_1,d-1} \vartheta_{1,d}^{\mu_1,d} \cdots \vartheta_{v_{d-1},d}^{\mu_1,d} \in (\bar{K}^*)^{G_d}.$$

Then based on the Remark 3 it follows that $\frac{\tilde{y}_{1,d}}{\sigma^{-1}(\tilde{y}_{1,d})} = \bar{y}_{j,d-1} \cdots \bar{y}_{j,d-1} \bar{y}_{j,d-1} \cdots \bar{y}_{j,d-1} \bar{y}_{j,d-1} \cdots \bar{y}_{j,d-1}$ and thus

$$\frac{g_{\ell}}{\sigma^{-1}(g_{\ell})} = \vartheta_{1,d-1}^{\mu_1,d-1} \cdots \vartheta_{v_{d-1},d-1}^{\mu_1,d-1} \vartheta_{1,d}^{\mu_1,d} \cdots \vartheta_{v_{d-1},d}^{\mu_1,d} \vartheta_{1,d}^{\mu_1,d} \cdots \vartheta_{v_{d-1},d}^{\mu_1,d} \vartheta_{1,d}^{\mu_1,d} \cdots \vartheta_{v_{d-1},d}^{\mu_1,d} = \vartheta_{d-1}^{\mu_1,d-1} \vartheta_{d-1}^{\mu_1,d} \vartheta_{d-1}^{\mu_1,d} \cdots \vartheta_{d-1}^{\mu_1,d} \vartheta_{d-1}^{\mu_1,d} \vartheta_{d-1}^{\mu_1,d} \cdots \vartheta_{d-1}^{\mu_1,d} \vartheta_{d-1}^{\mu_1,d} \cdots \vartheta_{d-1}^{\mu_1,d} = \rho_{d-1}(y_{d-1,d}) \rho_{d-1}(y_{d-1,d}) \rho_{d-1}(y_{d-1,d}) \cdots \rho_{d-1}(y_{d-1,d}) \rho_{d-1}(y_{d-1,d}) \rho_{d-1}(y_{d-1,d}) \cdots \rho_{d-1}(y_{d-1,d}) \rho_{d-1}(y_{d-1,d}) \cdots \rho_{d-1}(y_{d-1,d})$$

Hence also in this case we get

$$\frac{\sigma(g_{\ell})}{g_{\ell}} = \sigma(\frac{g_{\ell}}{\sigma^{-1}(g_{\ell})}) = \sigma\rho_{d-1}(y_{d-1,d}) = \rho_{d-1}(\alpha_{\ell,d}).$$
By iterative application of Lemma 8, the difference ring homomorphism \( \rho_{d-1} : \mathbb{A}_{d-1} \rightarrow \mathbb{G}_{d-1} \) can be extended to

\[
\rho_d : \mathbb{A}_{d-1} \langle y_{1,d} \rangle \langle y_{2,d} \rangle \cdots \langle y_{w_d,d} \rangle \rightarrow \mathbb{G}_{d-1} \langle \theta_{1,d} \rangle \langle \theta_{2,d} \rangle \cdots \langle \theta_{v_d,d} \rangle \langle \bar{y}_{1,d} \rangle \langle \bar{y}_{2,d} \rangle \cdots \langle \bar{y}_{c_d,d} \rangle
\]

with \( \rho_d|_{A_{d-1}} = \rho_{d-1} \) and \( \rho_d(y_{i,d}, d) = \bar{y}_i \) for \( 1 \leq i \leq w_d \). Finally, we show that for all \( f \in \mathbb{A}_d \) and \( n \in \mathbb{Z}_{\geq 0} \) we have \( \text{ev}(f, n) = \text{ev}(\rho_d(f), n) \). First note that for all \( n \geq 0 \) we have

\[
\text{ev}(y_{\ell,d}, n + 1) = \text{ev}(\sigma(y_{\ell,d}), n) = \text{ev}(\alpha_{\ell,d}, n) \text{ev}(y_{\ell,d}, n).
\]

(82)

On the other hand, since \( \rho_d \) is a difference ring homomorphism, we have that

\[
\sigma(\rho_d(y_{\ell,d})) = \rho_d(\sigma(y_{\ell,d})) = \rho_d(\alpha_{\ell,d}) \rho_d(y_{\ell,d}) = \rho_{d-1}(\alpha_{\ell,d}) \rho_d(y_{\ell,d})
\]

(83)

for all \( n \geq 0 \). Thus we get

\[
\text{ev}(\rho_d(y_{\ell,d}), n + 1) = \text{ev}(\sigma(\rho_d(y_{\ell,d})), n) = \text{ev}(\rho_{d-1}(\alpha_{\ell,d}), n) \text{ev}(\rho_d(y_{\ell,d}), n).
\]

(84)

By the induction hypothesis, \( \text{ev}(\alpha_{\ell,d}, n) = \text{ev}(\rho_{d-1}(\alpha_{\ell,d}), n) \) holds for all \( n \in \mathbb{Z}_{\geq 0} \). Therefore with (82) and (84) it follows that \( \text{ev}(y_{\ell,d}, n) \) and \( \text{ev}(\rho_d(y_{\ell,d}), n) \) satisfy the same first-order recurrence relation. With \( \text{ev}(y_{\ell,d}, 0) = 1 \) and

\[
\text{ev}(\rho_d(y_{\ell,d}, 0) = \text{ev}(\rho_d(y_{\ell,d}, 0) = \text{ev}(\theta_{1,d}, 0) \cdots \text{ev}(\theta_{v_d,d}, 0) \text{ev}(\bar{y}_{1,d}, 0) \cdots \text{ev}(\bar{y}_{c_d,d}, 0) = 1
\]

it follows then that \( \text{ev}(y_{\ell,d}, n) = \text{ev}(\rho_d(y_{\ell,d}, n) \) holds for all \( n \geq 0 \). Together with the induction hypothesis \( \text{ev}(f, n) = \text{ev}(\rho_d(f), n) \) for all \( f \in \mathbb{A}_{d-1} \) and \( n \in \mathbb{Z}_{\geq 0} \) we get (79) for all \( f \in \mathbb{A}_d \) and for all \( n \geq 0 \). Consequently also statement (1) of the Lemma holds.

Finally, if \( K \) is strongly \( \sigma \)-computable, the base case \( d = 1 \) can be executed explicitly by activating Lemma 7. In particular the induction step can be performed iteratively and thus the difference ring \( (\mathbb{G}_d, \sigma) \) with (73) together with (74), (75) and (76) can be computed. In addition, the difference ring \( (\mathbb{G}_d, \sigma) \), the difference ring homomorphism \( \rho_d : \mathbb{A}_d \rightarrow \mathbb{G}_d \) defined by (77) and the evaluation function \( \text{ev} \) can be computed. This completes the proof.

\[ \square \]

Remark 5. Note that the generators of \( \mathbb{G}_d \) with (73) constructed in Lemma 9 can be rearranged to get the AII-extension \( (\bar{K}[y_{1,1}] \cdots [y_{v_d,1}] \cdots [y_{1,d}] \cdots [y_{v_d,d}] [y_{1,1}] \cdots [y_{v_d,1}] y_{1,1} \cdots [y_{v_d,d}] d, \sigma) \) of \( (K, \sigma) \). Furthermore, a consequence of statement (3) of Lemma 9 is that the diagram

\[
\begin{array}{ccc}
\mathbb{A} & \xrightarrow{\psi} & S(K) \\
\rho \downarrow & & \downarrow \rho' \\
\mathbb{G}_d & \xrightarrow{\tilde{\tau}} & S(\bar{K})
\end{array}
\]

commutes where \( \mathbb{A} = \mathbb{A}_d \), \( \rho = \rho_d, \rho' = \text{id} \) and the difference ring homomorphism \( \tilde{\tau} \) and \( \psi \) are defined by \( \tilde{\tau}(f) = \langle \text{ev}(f, n) \rangle_{n \geq 0} \) and \( \psi(g) = \langle \text{ev}(g, n) \rangle_{n \geq 0} \) respectively.

Example 9 (Cont. Example 7). Take the ordered multiple chain AP-extension \( (A', \sigma) \) of \( (K, \sigma) \) with monomial depth \( 2 \) with \( A' = K[y_{1,1}] y_{1,1} y_{1,2} y_{2,1} y_{2,2} y_{3,1} y_{3,2} \) where \( (A', \sigma) \) is composed by the single chain AP-extensions of \( (K, \sigma) \) constructed in parts (41), (42), (43), (44), (45) and (46) of Example 7. By Lemma 9 and Remark 5 we can construct the AP-extension \( (G, \sigma) \) of \( (K, \sigma) \) where

\[
G = K[y_{1,1}] [\theta_{1,1}] [\theta_{1,2}] [\bar{y}_{1,1}] [\bar{y}_{1,2}] [\bar{y}_{3,1}] [\bar{y}_{3,2}]
\]

(85)

with the automorphism \( \sigma \) and evaluation function \( \text{ev} : G \times \mathbb{Z}_{\geq 0} \rightarrow K \) given by (47), (48) and

\[
\sigma(y_{1,1}) = 3 \sqrt{3} y_{1,1}, \quad \sigma(y_{2,1}) = 2 y_{2,1}, \quad \sigma(y_{3,1}) = 5 y_{3,1}, \quad \sigma(y_{2,2}) = 2 y_{2,1} y_{2,2}, \quad \sigma(y_{3,2}) = 5 y_{3,1} y_{3,2},
\]

\[
ev(y_{1,1}, n) = \prod_{k=1}^n \sqrt{3}, \quad \ev(y_{2,1}, n) = \prod_{k=1}^n 2, \quad \ev(y_{3,1}, n) = \prod_{k=1}^n 5, \quad \ev(y_{2,2}, n) = \prod_{k=1}^n \prod_{l=1}^k 2, \quad \ev(y_{3,2}, n) = \prod_{k=1}^n \prod_{l=1}^k 5
\]

(86)
with the following properties. By part [2] of Lemma 9, the sub-difference ring \( \mathcal{D} \) of the difference ring \( (\mathbb{G}, \sigma) \) with \( \mathbb{D} = \mathbb{K}(\hat{y}_{1,1}, \hat{y}_{2,1}, \hat{y}_{3,1}, \hat{y}_{2,2}, \hat{y}_{3,2}) \) is an ordered multiple chain II-extension of \((\mathbb{K}, \sigma)\) with the automorphism \( \sigma \) and the evaluation function \( e\hat{v} : \mathbb{D} \times \mathbb{Z}_{\geq 0} \to \mathbb{K} \) defined in [86]. In addition by part [3] of Lemma 9 we get the difference ring homomorphism \( \rho : \mathcal{A} \to \mathbb{G} \) defined by \( \rho|_{\mathbb{K}[[\hat{v}_{1,1}],[\hat{v}_{1,2}]]} = \text{id}_{\mathbb{K}[[\hat{v}_{1,1}],[\hat{v}_{1,2}]]} \) and

\[
\begin{align*}
\rho(y_{1,1}) &= \tilde{y}_{1,1}, & \rho(y_{2,1}) &= \tilde{y}_{2,1}, & \rho(y_{3,1}) &= \tilde{y}_{1,1}, & \rho(y_{4,1}) &= \tilde{y}_{3,1}, \\
\rho(y_{5,1}) &= \tilde{y}_{2,2}, & \rho(y_{2,2}) &= \tilde{y}_{2,2}, & \rho(y_{4,2}) &= \tilde{y}_{3,2}
\end{align*}
\]

such that \( e\hat{v}(\rho(f), n) = ev(f, n) \) holds for all \( n \in \mathbb{Z}_{\geq 0} \) and \( f \in \mathcal{A} \).

\section{Nested products with roots of unity}

Throughout this Subsection, \( \mathbb{K} \) is a field containing \( \mathbb{Q} \), \( \mathbb{K}_m \) is a splitting field for the polynomial \( x^m - 1 \) over \( \mathbb{K} \) (i.e., all roots of the polynomial \( x^m - 1 \) are in \( \mathbb{K}_m \)) for some \( m \in \mathbb{Z}_{\geq 2} \) and \( U_m \) is the set of all \( m \)-th roots of unity over \( \mathbb{K} \) in \( \mathbb{K}_m \) (which forms a multiplicative subgroup of \( \mathbb{K}^{\ast} \)). Then \( \text{Prod}_n(U_m) \) is the set of all geometric products over roots of unity in \( U_m \). For \( \mathbb{G} \subseteq \mathbb{K}_m(x) \) we define \( \text{Prod}_n(G; U_m) \) as the set of all elements

\[
\sum_{\mathbf{v}=(v_1,\ldots,v_e)\in S} a_{\mathbf{v}}(n) P_1(n)^{v_1} \cdots P_e(n)^{v_e}
\]

with \( e \in \mathbb{Z}_{\geq 0} \), \( S \subseteq \mathbb{Z}_{\geq 0}^e \) finite, \( a_{\mathbf{v}}(x) \in \mathbb{G} \) for \( \mathbf{v} \in S \) and \( P_1(n), \ldots, P_e(n) \in \text{Prod}_n(U_m) \).

The main result of this Subsection in Theorem 6 states that products over roots of unity with finite nesting depth can be represented by the single product \( \zeta^n \) where \( \zeta \in U_\lambda \) for some \( \lambda \geq 2 \) is a primitive \( \lambda \)-th root of unity.

\textbf{Theorem 6.} Suppose we are given the geometric products \( A_1(n), \ldots, A_e(n) \in \text{Prod}_n(U_m) \) in \( n \) of nesting depth \( r_i \in \mathbb{Z}_{\geq 0} \) with

\[
A_i(n) = \prod_{k_1=\ell_1}^{n} \zeta_{i,1}^{k_1} \prod_{k_2=\ell_2}^{n} \zeta_{i,2}^{k_2} \cdots \prod_{k_{r_i}=\ell_{r_i}}^{n} \zeta_{i,r_i}^{k_{r_i}}
\]

for \( 1 \leq i \leq e \) where \( \zeta_{i,j} \in U_m \), \( \ell_{i,j} \in \mathbb{Z}_{\geq 0} \) for \( 1 \leq j \leq r_i \). Then there exist a \( \lambda \in \mathbb{Z}_{\geq 2} \) with \( m \mid \lambda \) and a primitive \( \lambda \)-th root of unity \( \zeta_\lambda \in \mathbb{K}_\lambda \) satisfying the following property. For all \( 1 \leq i \leq e \) there exist \( f_{i,j} \in \mathbb{K}_\lambda \) for \( 0 \leq j < \lambda \) such that for

\[
B_i(n) = \sum_{j=0}^{\lambda-1} f_{i,j}(\zeta_\lambda^n)^j \in \text{Prod}_n(\mathbb{K}_\lambda; U_\lambda)
\]

we have

\[
A_i(n) = B_i(n)
\]

for all \( n \geq \max(\ell_1, \ldots, \ell_{r_i}) - 1 \). In particular, if \( \mathbb{K} \) is computable and one can solve Problem O (see below), the above construction can be given explicitly.

\subsection{The period and algorithmic aspects}

For the treatment of Theorem 6 we will introduce the period of a difference ring element introduced in [Kar81]. In particular, we will use the algorithms from [Sch10] that enable one to calculate the period within nested \( R \)-extensions, resp. \( A \)-extensions.

\textbf{Definition 15.} Let \( (\mathcal{A}, \sigma) \) be a difference ring. The \textit{period} of \( \alpha \in \mathcal{A}^\ast \) is defined by

\[
\text{per}(\alpha) = \begin{cases} 
0 & \text{if } \# n > 0 \text{ s.t. } \sigma^n(\alpha) = \alpha \\
\min\{n > 0 \mid \sigma^n(\alpha) = \alpha \} & \text{otherwise}
\end{cases}
\]
As it turns out, this task is connected to compute the order of a ring.

**Definition 16.** Let $A$ be a ring and let $\alpha \in A \setminus \{0\}$. Then the order of $\alpha$ is defined by

$$\text{ord}(\alpha) = \begin{cases} 0, & \text{if } \nexists n > 0 \text{ with } \alpha^n = 1. \\ \min\{n > 0 | \alpha^n = 1\}, & \text{otherwise}. \end{cases}$$

Namely, if we can solve the following problem (which is Problem GO with $w = 1$):

**Problem O for $\alpha \in K^*$**

*Given a field $K$ and $\alpha \in K^*$. Find ord($\alpha$).*

then we can also compute the period by the following lemma.

**Lemma 10.** Let $(E, \sigma)$ with $E = K_m[\vartheta_1] \ldots [\vartheta_e]$ be a simple $\mathbb{A}$-extension of $(K_m, \sigma)$. Then the following statements hold.

1. $\text{per}(\vartheta_i) > 0$ for all $1 \leq i \leq e$.
2. If $K_m$ is computable and Problem O is solvable, then $\text{per}(\vartheta_i)$ is computable for all $1 \leq i \leq e$.

**Proof.** Statement (1) follows by [Sch16, Proposition 5.5] (compare [Oca19, Proposition 6.2.20]) and statement (2) follows by [Sch16, Corollary 5.6] (compare [Oca19, Corollary 6.2.21]).

**Example 10 (Cont. Example 9).** Consider the sub-difference ring $(K[\vartheta_{1,1}][\vartheta_{1,2}], \sigma)$ of $(G, \sigma)$ with $-1$ and the evaluation function defined in (47) and (48). By construction it is a simple $\mathbb{A}$-extension of the difference field $(K, \sigma)$. Since $\text{per}(c) = 1$ for all $c \in K$, it follows that $\text{per}(c) = 1$. Applying the algorithms from [Sch16] (see the comments in the proof of Lemma 10) we compute for the depth-1 $\mathbb{A}$-monomial $\vartheta_{1,1}$ the period $\text{per}(\vartheta_{1,1}) = 2$, while for the depth-2 $\mathbb{A}$-monomial $\vartheta_{1,2}$ we get $\text{per}(\vartheta_{1,2}) = 4$.

### 5.3.2. Idempotent representation of single RP-extensions

In order to prove Theorem 6 we rely on the property that elements in basic RP-extensions can be expressed by idempotent elements. We start with the following basic facts inspired by [PS97].

**Lemma 11.** Let $F$ be a field and let $\zeta$ be a primitive $\lambda$-th root of unity. Let $F[\vartheta]$ be a polynomial ring subject to the relation $\vartheta^\lambda = 1$. Then the following statements hold.

1. The elements $e_0, \ldots, e_{\lambda - 1} \in F[\vartheta]$ with

$$e_k = e_k(\vartheta) := \prod_{i=0}^{\lambda-1} \frac{\vartheta - \zeta^i}{\zeta^{\lambda-1-k} - \zeta^i}$$

are idempotent and for all $0 \leq k < \lambda$ we have

$$e_k(\zeta^j) = \begin{cases} 1 & \text{if } j = \lambda - 1 - k \\ 0 & \text{if } j \neq \lambda - 1 - k \end{cases} \quad \text{and} \quad e_k(\vartheta^k) = e_{k+1} \mod \lambda.$$ (92)

2. The idempotent elements defined in (91) are pairwise orthogonal and $e_0 + \cdots + e_{\lambda - 1} = 1$. 

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In Proposition 3 we state that a simple RP-extension \((\mathbb{E}, \sigma)\) of a difference field \((\mathbb{F}, \sigma)\) with \(\mathbb{E} = \mathbb{F}[\vartheta](t_1 \ldots t_e)\) can be decomposed in terms of these idempotent elements. For more details in the general setting of RiIΣ-extensions we refer to [Sch17, Theorem 4.3]; compare also [PS97, Corollary 1.16], and [HS08, Lemma 6.8].

Proposition 3. Let \((\mathbb{E}, \sigma)\) with \(\mathbb{E} = \mathbb{F}[\vartheta](t_1 \ldots t_e)\) be an RP-extension of a difference field \((\mathbb{F}, \sigma)\) where \(\vartheta\) is an R-monomial of order \(\lambda\) with \(\zeta = \frac{\sigma(\vartheta)}{\vartheta} \in \text{const}(\mathbb{F}, \sigma)^*\) and the \(t_i\) are P-monomials. Let \(e_0, \ldots, e_{\lambda - 1}\) be the idempotent, pairwise orthogonal elements in \([91]\) that sum up to one. Then the following statements hold:

1. The ring \(\mathbb{E}\) can be written as the direct sum

\[
\mathbb{E} = e_0 \mathbb{E} \oplus \cdots \oplus e_{\lambda - 1} \mathbb{E}
\]

where \(e_k \mathbb{E}\) forms for all \(0 \leq k < \lambda\) a ring with \(e_k\) being the multiplicative identity element.

2. We have that \(e_k \mathbb{E} = e_k \tilde{\mathbb{E}}\) for \(0 \leq k < \lambda\) where \(\tilde{\mathbb{E}} = \mathbb{F}(t_1 \ldots t_e)\).

We are now ready to obtain the following key result; for the corresponding result for nested A-extensions of monomial depth 1 we refer to [Sch17, Lemma 2.22].

Theorem 7. Let \(m \in \mathbb{Z}_{\geq 2}\) and take a primitive \(m\)-th root of unity \(\zeta_m \in \mathbb{K}^*_m\). Let \((\mathbb{K}_m[\vartheta_1] \ldots [\vartheta_e], \sigma)\) be a simple A-extension of \((\mathbb{K}_m, \sigma)\) with \(\sigma(\vartheta_i) = \alpha_i \vartheta_i\) for \(1 \leq i \leq e\) where \(\alpha_i = \zeta_m^{u_i} \vartheta_1^{z_{i1}} \cdots \vartheta_e^{z_{ie}}\) with \(u_i, z_{ij} \in \mathbb{Z}_{\geq 0}\). Furthermore, let \(e_{\mathbb{K}_m}[\vartheta_1] \ldots [\vartheta_e] \times \mathbb{Z}_{\geq 0} \rightarrow \mathbb{K}_m\) be the evaluation function defined by

\[
e_{\mathbb{K}_m}(\vartheta_i, n) = \prod_{j=1}^{n} e_{\mathbb{K}_m}(\alpha_i, j - 1),
\]

and let \(\tau_m : \mathbb{K}_m[\vartheta_1] \ldots [\vartheta_e] \rightarrow \mathcal{S}(\mathbb{K})\) be the \(\mathbb{K}_m\)-homomorphism given by \(\tau_m(f) = \langle e_{\mathbb{K}_m}(f, n) \rangle_{n \geq 0}\). Then the following statements hold:

1. Define \(\lambda := \text{lcm}(\text{per}(\vartheta_1), \ldots, \text{per}(\vartheta_e)) > 1\). Then there is an R-extension \((\mathbb{K}_\lambda[\vartheta], \sigma)\) of \((\mathbb{K}_\lambda, \sigma)\) of order\(^{12}\) \(\lambda\) with \(\zeta = \frac{\sigma(\vartheta)}{\vartheta} \in \mathbb{K}_\lambda^*\) such that

\[
\varphi : \mathbb{K}_m[\vartheta_1] \ldots [\vartheta_e] \rightarrow \mathbb{K}_\lambda[\vartheta] = e_0 \mathbb{K}_\lambda \oplus \cdots \oplus e_{\lambda - 1} \mathbb{K}_\lambda
\]

defined with

\[
\varphi(f) = f_0 e_0 + \cdots + f_{\lambda - 1} e_{\lambda - 1}
\]

where \(f_i = e_{\mathbb{K}_m}(f, \lambda - 1 - i) \in \mathbb{K}_m \subseteq \mathbb{K}_\lambda\) for \(0 \leq i < \lambda\) is a difference ring homomorphism; here the \(e_k\) are the idempotent orthogonal elements defined in \([91]\). In particular, \(\varphi|_{\mathbb{K}_m} = \text{id}_{\mathbb{K}_m}\).

2. Take the evaluation function \(e_{\mathbb{K}_\lambda}[\vartheta] \times \mathbb{Z}_{\geq 0} \rightarrow \mathbb{K}_\lambda\) defined by \(e_{\mathbb{K}_\lambda}[\vartheta]|_{\mathbb{K}_\lambda} = \text{id}\) and \(e_{\mathbb{K}_\lambda}(\vartheta, n) = \zeta^n\) and consider the \(\mathbb{K}_\lambda\)-homomorphism \(\tau_\lambda : \mathbb{K}_\lambda[\vartheta] \rightarrow \mathcal{S}(\mathbb{K}_\lambda)\) defined by \(\tau_\lambda(f) = \langle e_{\mathbb{K}_\lambda}(f, n) \rangle_{n \geq 0}\). Then for the pairwise orthogonal elements \(e_k\) defined in \([91]\) with \(0 \leq k < \lambda\), we have that

\[
e_{\mathbb{K}_\lambda}(e_k, n) = \begin{cases} 1 & \text{if } \lambda | n + k + 1, \\ 0 & \text{if } \lambda \nmid n + k + 1. \end{cases}
\]

3. The \(\mathbb{K}_\lambda\)-homomorphism \(\tau_\lambda : \mathbb{K}_\lambda[\vartheta] \rightarrow \mathcal{S}(\mathbb{K}_\lambda)\) with the evaluation function defined in part (2) is injective.

\(^{12}\)\(\mathbb{K}_\lambda\) is a finite algebraic extension of \(\mathbb{K}_m\) and \(\zeta \in \mathbb{K}_m\) is a primitive \(\lambda\)-th root of unity.
The diagram
\[
\begin{array}{c}
\mathbb{K}_m[\vartheta_1] \cdots [\vartheta_e] \xrightarrow{\tau_m} S(\mathbb{K}_m) \\
\varphi \downarrow \\
\mathbb{K}_\lambda[\vartheta] \simeq e_0 \mathbb{K}_\lambda \oplus \cdots \oplus e_{\lambda-1} \mathbb{K}_\lambda \xrightarrow{\tau_\lambda} S(\mathbb{K}_\lambda)
\end{array}
\]
(98)

commutes where \( \varphi' : S(\mathbb{K}_m) \rightarrow S(\mathbb{K}_\lambda) \) is the injective difference ring homomorphism defined by \( \varphi'(a) = a \).

If \( \mathbb{K}_m \) is computable and Problem \( Q \) is solvable in \( \mathbb{K}_m \), then the above constructions are computable.

**Proof.**

(1) Since \( c_m^\sigma \in \mathbb{K}_m^* \), per(\( c_m^\sigma \)) = 1 > 0 for all 1 \( \leq \) \( u_i \leq \) \( e \). In addition, it follows by statement \([1]\) of Lemma \([10]\) that per(\( \vartheta_i \)) > 0 for all 1 \( \leq \) \( i \leq \) \( e \). Define \( \lambda := \text{lcm}(m, \text{per}(\vartheta_1), \ldots, \text{per}(\vartheta_e)) > 1 \). Note that \( m \mid \lambda \), i.e., \( \mathbb{K}_\lambda \) is an algebraic field extension of \( \mathbb{K}_m \). Finally, take \( \zeta := e^{\frac{\lambda}{m}} \in \mathbb{K}_\lambda \)
and construct the A-extension (\( \mathbb{K}_\lambda[\vartheta], \sigma \)) of (\( \mathbb{K}_\lambda, \sigma \)) with \( \sigma(\vartheta) = \zeta \vartheta \). By Lemma \([3]\) it follows that (\( \mathbb{K}_\lambda[\vartheta], \sigma \)) is an R-extension of (\( \mathbb{K}_\lambda, \sigma \)). By Proposition \([3]\) we have that \( \mathbb{K}_\lambda[\vartheta] = e_0 \mathbb{K}_\lambda \oplus \cdots \oplus e_{\lambda-1} \mathbb{K}_\lambda \)
where the \( e_k \) for \( 0 \leq k < \lambda \) are the orthogonal idempotent elements defined in \([9]\). Now consider the map \([55]\) defined by \([56]\). We will now show that \( \varphi \) is a ring homomorphism. Observe that for any \( c \in \mathbb{K}_m \), \( ev_m(c, i) = c \) for all \( i \in \mathbb{Z}_{\geq 0} \) and with statement \([2]\) of Lemma \([1]\) we have that
\[
\varphi(c) = c e_0 + \cdots + c e_{\lambda-1} = c(e_0 + \cdots + e_{\lambda-1}) = c.
\]
Further, let \( f, g \in \mathbb{K}_m[\vartheta_1] \cdots [\vartheta_e] \) with \( f := a \vartheta_1^{\sigma_1} \cdots \vartheta_e^{\sigma_e} \) and \( g := b \vartheta_1^{\hat{\sigma}_1} \cdots \vartheta_e^{\hat{\sigma}_e} \) where \( a, b \in \mathbb{K}_m \) and \( \nu_i, \lambda_i \in \mathbb{Z}_{\geq 0} \) for \( 1 \leq i \leq e \). Define \( f_k := ev_m(f, \lambda - 1 - k) \) and \( g_k := ev_m(g, \lambda - 1 - k) \) for \( 0 \leq k < \lambda \). Then,
\[
\varphi(f + g) = ev_m(f + g, \lambda - 1) e_0 + \cdots + ev_m(f + g, 0) e_{\lambda-1}
\]
\[
= (ev_m(f, \lambda - 1) + ev_m(g, \lambda - 1)) e_0 + \cdots + (ev_m(f, 0) + ev_m(g, 0)) e_{\lambda-1}
\]
\[
= (ev_m(f, \lambda - 1) e_0 + \cdots + ev_m(f, 0) e_{\lambda-1}) + (ev_m(g, \lambda - 1) e_0 + \cdots + ev_m(g, 0) e_{\lambda-1})
\]
\[
= (f_0 e_0 + \cdots + f_{\lambda-1} e_{\lambda-1}) + (g_0 e_0 + \cdots + g_{\lambda-1} e_{\lambda-1})
\]
\[
= \varphi(f) + \varphi(g).
\]
Similarly,
\[
\varphi(f g) = ev_m(f g, \lambda - 1) e_0 + \cdots + ev_m(f g, 0) e_{\lambda-1}
\]
\[
= (ev_m(f, \lambda - 1) ev_m(g, \lambda - 1)) e_0 + \cdots + (ev_m(f, 0) ev_m(g, 0)) e_{\lambda-1}
\]
\[
= f_0 g_0 e_0 + f_1 g_1 e_1 + \cdots + f_{\lambda-1} g_{\lambda-1} e_{\lambda-1}
\]
\[
= (f_0 e_0 + \cdots + f_{\lambda-1} e_{\lambda-1}) (g_0 e_0 + \cdots + g_{\lambda-1} e_{\lambda-1})
\]
\[
= \varphi(f) \varphi(g).
\]
The first equality follows since the \( e_i \) are idempotent. Thus, \( \varphi \) is a ring homomorphism. Next we show by induction on the number of A-monomials, \( e \in \mathbb{Z}_{\geq 0} \), that \( \varphi \) is a difference ring homomorphism. For the base case, i.e., \( e = 0 \), there are no A-monomials. Since \( \sigma(\varphi(c)) = \sigma(c) = c = \varphi(c) = \varphi(\sigma(c)) \) for all \( c \in \mathbb{K}_m \), \( \varphi \) is a difference ring-homomorphism. Now assume that the statement holds for all A-monomials \( \vartheta_i \) with \( 0 \leq i < e \), and consider an A-monomial \( \vartheta_e \) with \( \sigma(\vartheta_e) = \hat{\alpha} \vartheta_e \) where \( \hat{\alpha} \in (\mathbb{K}_m^*, \mathbb{K}_m[\vartheta_1] \cdots [\vartheta_{e-1}], \mathbb{K}_m) \). Then we will show that
\[
\sigma(\varphi(\vartheta_e)) = \varphi(\sigma(\vartheta_e))
\]
(99)
holds. For the left hand side of (99), we have that \( \varphi(\vartheta_e) = \gamma_0 e_0 + \cdots + \gamma_{\lambda-1} e_{\lambda-1} \) where \( \gamma_i = ev_m(\vartheta_e, \lambda - 1 - i) \in \mathbb{K}_m \) for \( 0 \leq i < \lambda \) are \( \lambda \)-th roots of unity. Thus,

\[
\sigma(\varphi(\vartheta_e)) = \sigma(\gamma_0) \sigma(e_0) + \cdots + \sigma(\gamma_{\lambda-1}) \sigma(e_{\lambda-1}).
\]

By (92) we have that \( \sigma(e_{\lambda-1}) = e_0 \) and \( \sigma(e_i) = e_{i+1} \) for \( 0 \leq i < \lambda - 1 \). In addition, for \( 1 \leq i < \lambda \) we get \( \sigma(\gamma_i) = ev_m(\vartheta_e, \lambda - i) = \gamma_i^{-1} \). For \( i = \lambda \) observe that per(\( \vartheta_e \)) \| \lambda \) by definition and thus \( \sigma^\lambda(\vartheta_e) = \vartheta_e \). Consequently \( \sigma(\gamma_0) = ev_m(\vartheta_e, \lambda) = ev_m(\sigma^\lambda(\vartheta_e), 0) = ev_m(\vartheta_e, 0) = \gamma_{\lambda-1} \). Therefore,

\[
\sigma(\varphi(\vartheta_e)) = \gamma_0 e_0 + \cdots + \gamma_{\lambda-1} e_{\lambda-1}
\]

where \( \gamma_0 = \gamma_{\lambda-1} \) and \( \gamma_i = \gamma_{\lambda-1}^{-1} \) for \( 1 \leq i < \lambda - 1 \).

For the right hand side of (99), we have

\[
\varphi(\sigma(\vartheta_e)) = \varphi(\tilde{\vartheta}_e) = \varphi(\tilde{\vartheta}_e) e_0 \varphi(\vartheta_e) = (\alpha_0 e_0 + \cdots + \alpha_{\lambda-1} e_{\lambda-1})(\gamma_0 e_0 + \cdots + \gamma_{\lambda-1} e_{\lambda-1})
\]

\[
= \alpha_0 \gamma_0 e_0 + \cdots + \alpha_{\lambda-1} \gamma_{\lambda-1} e_{\lambda-1}
\]

(101)

where \( \alpha_i = ev_m(\tilde{\vartheta}_e, \lambda - 1 - i) \) and \( \gamma_i = ev_m(\vartheta_e, \lambda - 1 - i) \) for \( 0 \leq i < \lambda \) are \( \lambda \)-th roots of unity. Again (101) holds since the \( e_i \) are idempotent. Finally observe that for \( 0 \leq i < \lambda \) we have

\[
\alpha_i \gamma_i = ev_m(\vartheta_e, \lambda - 1 - i) = ev_m(\tilde{\vartheta}_e, \lambda - 1 - i) = ev_m(\vartheta_e, \lambda - 1 - i) = ev_m(\sigma(\vartheta_e), \lambda - 1 - i) = ev_m(\vartheta_e, \lambda - i) = \gamma_i.
\]

With (100) we conclude that (99) holds. Thus, \( \varphi \) is a difference ring homomorphism.

(2) By Lemma [2] we can define the evaluation function \( ev_\lambda : \mathbb{K}_\lambda[\vartheta] \times \mathbb{Z}_{\geq 0} \rightarrow \mathbb{K}_\lambda \) with \( ev_\lambda|_\mathbb{K}_\lambda = \text{id} \) and \( ev_\lambda(\vartheta, n) = \zeta^n \) and by Lemma [4] we get the \( \mathbb{K}_\lambda \)-homomorphism \( \tau_\lambda : \mathbb{K}_\lambda[\vartheta] \rightarrow \mathcal{S}(\mathbb{K}_\lambda) \) defined by \( \tau_\lambda(\vartheta) = (ev_\lambda(f, n))_{n \geq 0} \). Statement (97) follows by (92).

(3) Since \( (\mathbb{K}_\lambda[\vartheta], \sigma) \) is an R-extension of a difference field \((\mathbb{K}_\lambda, \sigma)\) it follows by Theorem [2] that \( \tau_\lambda \) is injective.

(4) Let \( \alpha \in \mathbb{K}_m[\vartheta_1, \ldots, \vartheta_c] \) and let \( ev_m, ev_\lambda \) be evaluation functions for \( \mathbb{K}_m[\vartheta_1, \ldots, \vartheta_c] \) and \( \mathbb{K}_\lambda[\vartheta] \) defined by (94) and (97), respectively. We will show

\[
\varphi'(\tau_\lambda(\alpha)) = \tau_\lambda(\varphi(\alpha))
\]

(102)

For the left hand side of (102), we have

\[
\varphi'(\tau_\lambda(\alpha)) = \tau_\lambda(\alpha) = \langle ev_m(\alpha, n) \rangle_{n \geq 0} \in \mathcal{S}(\mathbb{K}_m) \subseteq \mathcal{S}(\mathbb{K}_\lambda).
\]

For the right hand side of (102) we note by (96) that \( \varphi(\alpha) = \alpha_0 e_0 + \cdots + \alpha_{\lambda-1} e_{\lambda-1} \) holds where \( \alpha_i = ev_m(\vartheta_e, \lambda - 1 - i) \in \mathbb{K}_m \subseteq \mathbb{K}_\lambda \) for \( 0 \leq i < \lambda \). Thus,

\[
\tau_\lambda(\varphi(\alpha)) = \langle ev_\lambda(\alpha_0 e_0 + \cdots + \alpha_{\lambda-1} e_{\lambda-1}, n) \rangle_{n \geq 0}
\]

\[
= \langle ev_\lambda(\alpha_0 e_0, n) \rangle_{n \geq 0} + \cdots + \langle ev_\lambda(\alpha_{\lambda-1} e_{\lambda-1}, n) \rangle_{n \geq 0}
\]

\[
= \alpha_0 \langle ev_\lambda(e_0, n) \rangle_{n \geq 0} + \cdots + \alpha_{\lambda-1} \langle ev_\lambda(e_{\lambda-1}, n) \rangle_{n \geq 0}
\]

\[
= \langle ev_m(\alpha, n) \rangle_{n \geq 0}.
\]

The last equality follows by (97). This implies that the diagram (98) commutes.
Finally, if $K_m$ is computable and Problem O is solvable in $K_m$, then by statement (2) of Lemma 10, $\text{per}(\vartheta_i)$ is computable for all $1 \leq i \leq e$. Consequently, the R-extension $(K_\lambda[\vartheta], \sigma)$ of $(K_\lambda, \sigma)$, the evaluation function $\text{ev}_\lambda : K_\lambda[\vartheta] \times \mathbb{Z}_{\geq 0} \rightarrow K_\lambda$ given in statement (2) and the injective $K_\lambda$-homomorphism $\tau_\lambda : K_\lambda[\vartheta] \rightarrow S(K_\lambda)$ given in statement (4) can be constructed explicitly. \hfill $\square$

**Remark 6.** By statement (4) of Theorem 7 and Remark 7 we observe that for a fixed $k \in \mathbb{Z}_{\geq 0}$ and $\alpha \in K_m[\vartheta_1] \ldots [\vartheta_e]$ we get

\[
\text{ev}_m(\alpha, k) = \text{ev}_\lambda(\varphi(\alpha), k) = \alpha_0 \text{ev}_\lambda(e_0, k) + \cdots + \alpha_{\lambda-1} \text{ev}_\lambda(e_{\lambda-1}, k) = \alpha_j \text{ev}_\lambda(e_j, k) = \alpha_j = \text{ev}_m(\alpha, j)
\]

for some $j \in \{0, 1, \ldots, \lambda - 1\}$ with $\lambda | k - j$. In other words, the sequence repeats periodically.

**Example 11 (Cont. Example 10).** Consider the $U_2$-simple A-extension $(K[\vartheta_1, 1][\vartheta_1, 2], \sigma)$ of $(K, \sigma)$ with the automorphism and the evaluation function given in (47) and (48), which was constructed in Examples 9 and 10 with $K = \mathbb{Q}(\sqrt{3})(= K_2)$. From Example 10, we already know the period of the A-monomials $\vartheta_1, 1$ and $\vartheta_1, 2$ in $K[\vartheta_1, 1][\vartheta_1, 2]$. Set $\lambda = \text{lcm}(m, \text{per}(\vartheta_1, 1), \text{per}(\vartheta_1, 2)) = 4$ with $m = 2$, take a primitive 4th root of unity, say $\zeta := e^{\frac{2\pi i}{4}} = (-1)^{\frac{1}{2}} = i$ and define $\tilde{K} = \mathbb{Q}(i, \sqrt{3})(= K_4)$. Then by statement (1) of Theorem 7, there is an R-extension $(\tilde{K}[\vartheta], \sigma)$ of $(\tilde{K}, \sigma)$ of order 4 with the automorphism

\[
\sigma(\vartheta) = \vartheta
\]

and the evaluation function $\tilde{\text{ev}} : \tilde{K}[\vartheta] \times \mathbb{Z}_{\geq 0} \rightarrow \tilde{K}$ given by

\[
\tilde{\text{ev}}(\vartheta, n) = \prod_{k=1}^{n} \vartheta = i^n.
\]

We have $\tilde{K}[\vartheta] = e_0 \tilde{K} + e_1 \tilde{K} + e_2 \tilde{K} + e_3 \tilde{K}$ where the idempotent elements $e_k$ for $0 \leq k \leq 3$ are defined by

\[
e_0 = \frac{1}{4} \left( \vartheta^3 + i \vartheta^2 - \vartheta - i \right), \quad e_1 = \frac{1}{4} \left( 1 - \vartheta + \vartheta^2 - \vartheta^3 \right),
\]

\[
e_2 = \frac{1}{4} \left( -\vartheta^3 + i \vartheta^2 + \vartheta - i \right), \quad e_3 = \frac{1}{4} \left( 1 + \vartheta + \vartheta^2 + \vartheta^3 \right)
\]

with $e_0 + e_1 + e_2 + e_3 = 1$. Furthermore, the ring homomorphism $\varphi : K[\vartheta_1, 1][\vartheta_1, 2] \rightarrow \tilde{K}[\vartheta]$ defined by $\varphi[\vartheta] = \text{id}_{K}$ and $\varphi(\vartheta_1, 1) = \beta_{0,0} e_0 + \beta_{1,1} e_1 + \beta_{2,2} e_2 + \beta_{3,3} e_3$ where $\beta_{i,j} = \text{ev}(\vartheta_{i,j}, 3 - j)$ for $i \in \{1, 2\}$ and $0 \leq j \leq 3$ is a difference ring homomorphism. More precisely, for the A-monomials we have that

\[
\varphi(\vartheta_1, 1) = -e_0 + e_1 - e_2 + e_3 = \vartheta^2,
\]

\[
\varphi(\vartheta_1, 2) = -e_0 - e_1 + e_2 + e_3 = \frac{(1-i)}{2} \vartheta(\vartheta^2 + i).
\]

Given $\text{ev}$ and $\tilde{\text{ev}}$ we obtain the difference ring homomorphisms $\tau_2 : K[\vartheta_1, 1][\vartheta_1, 2] \rightarrow S(K)$ defined by $\tau_2(f) = (\text{ev}(f, n))_{n \geq 0}$ and $\tau_4 : \tilde{K}[\vartheta] \rightarrow S(\tilde{K})$ defined by $\tau_4(f) = (\text{ev}(f, n))_{n \geq 0}$. In particular, by statement (3) of Theorem 7, $\tau_4$ is injective. Finally, by defining the embedding $\varphi' : S(\tilde{K}) \rightarrow S(K)$ with $\varphi'(a) = a$ for all $a \in S(\tilde{K})$ we conclude by statement (4) of the Theorem 7 that the following diagram commutes

\[
\begin{array}{ccc}
K[\vartheta_1, 1][\vartheta_1, 2] & \xrightarrow{\tau_2} & S(\tilde{K}) \\
\varphi \downarrow & & \downarrow \varphi' \\
\tilde{K}[\vartheta] = e_0 \tilde{K} + e_1 \tilde{K} + e_2 \tilde{K} + e_3 \tilde{K} & \xrightarrow{\tau_4} & S(\tilde{K}).
\end{array}
\]

We are finally ready to obtain the
Proof of Theorem 6. Suppose we are given the geometric products $A_1(n), \ldots, A_e(n) \in \text{Prod}_n(U_m)$ in $n$ with (88) where $\zeta_i,r_i \neq 1$ for $1 \leq i \leq e$. As elaborated in Section 2.2 we can rewrite each $A_i(n)$ as

$$A_i(n) = u_i \hat{A}_i(n) \quad \text{where} \quad \hat{A}_i(n) = \prod_{k_1=1}^{n} \prod_{k_2=1}^{k_1} \prod_{k_{r_i}=1}^{k_{r_{i-1}}} \zeta_{i,r_i}$$

and $u_i \in U_m$ which holds for all $n \geq \max(\ell_{i,1}, \ldots, \ell_{i,r_i}) - 1 =: \delta_i$. Similar to Remark 4 we can rephrase the products in a simple $A$-extension $(\mathcal{A}, \sigma)$ of $(K_m, \sigma)$ with $\mathcal{A} = K_m[\vartheta_{1,1}, \ldots, \vartheta_{1,r_1}, \ldots, \vartheta_{e,1}, \ldots, \vartheta_{e,r_e}]$ where $\alpha_{i,j} := \sigma(\vartheta_{i,j}) = \zeta_{m,i} \vartheta_{i,1} \cdots \vartheta_{i,j-1}$ for $1 \leq i \leq e$ and $1 \leq j \leq r_i$ with $u_i \in U_m$ equipped with the evaluation function $ev_m : \mathcal{A} \times \mathbb{Z}_{\geq 0} \to K_m$ defined by $ev_m(\vartheta_{i,j}, n) = \prod_{k=1}^{n} ev_m(\alpha_{i,j}, k - 1)$ with the following property. For all $i$ with $1 \leq i \leq e$, there are $\nu_i, \mu_i$ such that the geometric product $\hat{A}_i(n)$ is modelled by $\vartheta_{\nu_i, \mu_i}$, i.e.,

$$ev_m(\vartheta_{\nu_i, \mu_i}, n) = \hat{A}_i(n)$$

holds for all $n \geq 0$. In particular, we get the $K$-homomorphism $\tau_m : \mathcal{A} \to S(K_\lambda)$. By Theorem 7 there is a single $R$-extension $(K_\lambda, \sigma)$ of $(K_m, \sigma)$ subject to the relations $\vartheta^k = 1$ and $\sigma(\vartheta) = \zeta_\lambda \vartheta$ where $\lambda := \lcm(m, \text{per}(\gamma_1), \ldots, \text{per}(\gamma_e)) > 0$, $\zeta_\lambda := e^{\frac{\pi i}{\lambda}} = (-1)^\frac{\lambda}{2} \in K_\lambda$ and $K_\lambda$ is some algebraic extension of $K_m$ with $m | \lambda$. Furthermore, there is a difference ring homomorphism $\phi : \mathcal{A} \to K_\lambda$ and a $K_\lambda$-embedding $\tau_\lambda : K_\lambda[\vartheta] \to S(K_\lambda)$ with $\tau_\lambda(\vartheta) = (\zeta_\lambda)^n$ such that $\tau_m(f) = \tau_\lambda(\phi(f))$ holds for all $f \in \mathcal{A}$. In particular, we get

$$ev_m(\vartheta_{\nu_i, \mu_i}, n) = ev_\lambda(\phi(\vartheta_{\nu_i, \mu_i}), n)$$

for $1 \leq i \leq e$. Now define $g_{i,k} \in K_\lambda$ by $\phi(\vartheta_{\nu_i, \mu_i}) = \sum_{k=0}^{n-1} g_{i,k} \vartheta^k \in K_\lambda$. Then for $G_i(n) := \sum_{k=0}^{n-1} g_{i,k} (\zeta_\lambda)^k$ with $1 \leq i \leq e$ we get

$$ev_m(\vartheta_{\nu_i, \mu_i}, n) \stackrel{(109)}{=} ev_\lambda(\phi(\vartheta_{\nu_i, \mu_i}), n) = G_i(n) \quad \forall n \geq 0.$$ 

With (108) and (109) we conclude that

$$A_i(n) = u_i \hat{A}_i(n) = u_i G_i(n)$$

holds for all $n \geq \delta_i$. In particular, for $B_i(n)$ given in (89) with $f_{i,k} := u_i g_{i,k} \in K_m$ we get (90).

If $K_m$ is computable and Problem 0 can be solved, Theorem 7 is constructive and all the above ingredients can be given explicitly. \hfill \Box

Example 12 (Cont. Example 11). Consider the product expression

$$A(n) = \sqrt{3} \prod_{i=1}^{n} (-1) + 2 \prod_{k=1}^{n} \prod_{i=1}^{k} (-1) + 3 \prod_{i=1}^{n} (-1) \prod_{k=1}^{n} \prod_{i=1}^{k} (-1) \in \text{Prod}_n \left( \mathbb{Q}(\sqrt{3}), U_2 \right).$$

For this instance we follow the construction in Example 11 and get $f = \sqrt{3} \vartheta_{1,1} + 2 \vartheta_{1,2} + 3 \vartheta_{1,1} \vartheta_{1,2} \in K[\vartheta_{1,1}][\vartheta_{1,2}]$ with $ev(f, n) = A(n)$ for all $n \in \mathbb{Z}_{\geq 0}$. As a consequence we obtain $f := \phi(f) = (\frac{1}{2} - \frac{1}{2} i) \vartheta ((2 + 3i) \vartheta^2 + (1 + i) \sqrt{3} \vartheta + (3 + 2i)) \in K[\vartheta]$ yielding for $n \in \mathbb{Z}_{\geq 0}$ the identity

$$A(n) = ev(f, n) = ev(\tilde{f}, n) = \left( \frac{1}{2} - \frac{1}{2} i \right) n \left( (2 + 3i)(i^n)^2 + (1 + i) \sqrt{3} n + (3 + 2i) \right).$$

In particular, as claimed in Theorem 8 each of the products in $A(n)$ can be expressed in terms of $i^n$. Namely, for all $n \in \mathbb{Z}_{\geq 0}$ we obtain

$$\prod_{k=1}^{n} (-1) = ev(\vartheta_{1,1}, n) = ev(\vartheta(\vartheta), n) = ev(\vartheta^2, n) = (i^n)^2,$n

$$\prod_{k=1}^{n} \prod_{i=1}^{k} (-1) = ev(\vartheta_{1,2}, n) = ev(\vartheta(\vartheta), n) = ev(\frac{1-i}{2} \vartheta (\vartheta^2 + i), n) = \frac{1-i}{2} \vartheta (\vartheta^2 + i).$$

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6. A complete solution of Problem RPE

We are now ready to combine the building blocks from the previous section to solve Problem RPE in Sections 6.1 and 6.2 below. Afterwards we apply in Section 6.3 the machinery implemented within the package Nested Products to concrete examples.

6.1. The difference ring setting for nested geometric products

First we combine Lemma 9 discussed in Subsection 5.2 and Theorem 7 discussed in Subsection 5.3. As a consequence, we will obtain the necessary difference ring tools for the full treatment of geometric products of arbitrary but finite nesting depth.

Theorem 8. For \(1 \leq \ell \leq m\), let \((\mathbb{K}_\ell, \sigma)\) with \(\mathbb{K}_\ell = \mathbb{K}(y_{\ell,1}) \cdots (y_{\ell,n})\) be the single chain \(P\)-extensions of \((\mathbb{K}, \sigma)\) over \(\mathbb{K} = \mathbb{K}(s_1, \ldots, s_m)\) with base \(h_\ell \in \mathbb{K}^*\) for \(1 \leq \ell \leq m\), the automorphisms (71) and the evaluation functions (72). Let \(d := \max(s_1, \ldots, s_m)\) and \(A_0 = \mathbb{K}\). Consider the tower of difference ring extensions \((A_i, \sigma)\) of \((A_{i-1}, \sigma)\) where \(A_i = A_{i-1}(y_{i,1})(y_{i,2}) \cdots (y_{i,n})\) for \(1 \leq i \leq d\) with \(m = w_1 \geq w_2 \geq \cdots \geq w_d\) and the automorphism (71) and the evaluation function (72). This yields the ordered multiple chain \(P\)-extension \((A_d, \sigma)\) of \((\mathbb{K}, \sigma)\) of monomial depth at most \(d\) composed by the single chain \(P\)-extensions \((\mathbb{K}_\ell, \sigma)\) of \((\mathbb{K}, \sigma)\) for \(1 \leq \ell \leq m\) with (71) and (72). Then one can construct

1. an \(R\Pi\)-extension \((\mathbb{D}, \sigma)\) of \((\mathbb{K}', \sigma)\) with
   \[
   \mathbb{D} = \mathbb{K}'[\vartheta](\tilde{y}_{1,1}) \cdots (\tilde{y}_{1,1}) \cdots (\tilde{y}_{d,d})
   \]
   where \(\mathbb{K}' = K'(\kappa_1, \ldots, \kappa_u)\) and \(K'\) is a finite algebraic field extension of \(K\) and with the automorphism
   \[
   \sigma(\vartheta) = \zeta' \vartheta \quad \text{and} \quad \sigma(\tilde{y}_{\ell,d}) = \tilde{\alpha}_{\ell,k} \tilde{y}_{\ell,k}
   \]
   where \(\zeta' \in K'\) is a \(\lambda\)'-th root of unity and
   \[
   \tilde{\alpha}_{\ell,k} = h_\ell \tilde{y}_{\ell,1} \cdots \tilde{y}_{\ell,k-1} \in (\mathbb{K}'(s))_{\mathbb{K}}(\tilde{y}_{\ell,1}) \cdots (\tilde{y}_{\ell,k-1})
   \]
   for \(1 \leq k \leq d\) and \(1 \leq \ell \leq e_k\);
2. an evaluation function \(\tilde{\text{ev}} : \mathbb{D} \times \mathbb{Z}_{\geq 0} \to \mathbb{K}'\) defined as
   \[
   \tilde{\text{ev}}(\vartheta, n) = \prod_{j=1}^{n} \zeta'\quad \text{and} \quad \tilde{\text{ev}}(\tilde{y}_{\ell,k}, n) = \prod_{j=1}^{n} \tilde{\text{ev}}(\tilde{\alpha}_{\ell,k}, j - 1);
   \]
3. a difference ring homomorphism \(\varphi : A_d \to \mathbb{D}\) defined by \(\varphi|_{\mathbb{K}} = \text{id}_{\mathbb{K}}\) and
   \[
   \varphi(y_{\ell,k}) = \gamma_{\ell,k} \tilde{y}_{\ell,1} \cdots \tilde{y}_{\ell,s_k}
   \]
   for \(1 \leq \ell \leq m\) and \(1 \leq k \leq s_\ell\) with \(\gamma_{\ell,k} \in \mathbb{K}'[\vartheta]\) and \(v_{\ell,i,k} \in \mathbb{Z}\) for \(1 \leq i \leq e_k\)

such that for all \(f \in A_d\) and \(n \in \mathbb{Z}_{\geq 0}\) we have

\[
\text{ev}(f, n) = \tilde{\text{ev}}(\varphi(f), n).
\]

If \(\mathbb{K}\) is strongly \(\sigma\)-computable, then the constructions above are computable.

---

13 For concrete instances the R-monomial \(\vartheta\) might not be needed. In particular, if \(v_{\ell,1,k} = 0\) in (112), it can be removed.
14 Note that for all \(c \in \mathbb{K}', \tilde{\text{ev}}(c, n) = c\) for all \(n \geq 0\).
Proof. Let \((A_d, \sigma)\) be an ordered multiple chain P-extension of \((K, \sigma)\) of monomial depth at most \(d\) with the automorphism \(\sigma : A_d \to A_d\) defined by \((72)\) and the evaluation function \(ev : A_d \times \mathbb{Z}_{\geq 0} \to K\) defined by \((71)\). Then by Lemma 9 we can construct an ordered multiple chain AP-extension \((G_d, \sigma)\) of \((K, \sigma)\) of monomial depth at most \(d\) with \(G = \tilde{K}(\kappa_1, \ldots, \kappa_u)\), \(\tilde{K}\) being a finite algebraic field extension of \(K\) where where \(G_d\) is given by \((70)\) with the automorphism \((74)\) and \((75)\), the evaluation function \(ev : \mathbb{Z}_{\geq 0} \to \tilde{K}\) where \(\tilde{ev}\) is replaced by \(ev\), and the difference ring homomorphism \(\rho_d : \tilde{A}_d \to \tilde{G}_d\) defined by \((77)\) with \(\rho_d|_K = id_K\) and \((78)\) with the following properties: the sub-difference ring \((\tilde{A}_d, \sigma)\) of \((\tilde{G}_d, \sigma)\) where \(\tilde{A}_d\) is given by \((78)\) is a \(\Pi\)-extension of \((K, \sigma)\). Furthermore, for all \(f \in \tilde{A}_d\) and for all \(n \in \mathbb{Z}_{\geq 0}\) we have

\[ev(f, n) = ev'(\rho_d(f), n)\]  

(144)

If \(v_1 = 0\) (and thus \(v_2 = \cdots = v_d = 0\)), i.e., no A-monomials are involved, we are essentially done. We simply adjoin a redundant R-monomial (compare the footnote in Theorem 8). Otherwise \(v_1 \geq 1\) and by Remark 3 the generators in \(G_d\) can be rearranged to get the AII-extension \((\tilde{G}, \sigma)\) of \((K, \sigma)\) where

\[\tilde{G} = \tilde{K}[\theta_{v_1,1}] \cdots [\theta_{v_d,1}] \cdots [\theta_{v_d,d}]\langle \tilde{g}_1,1 \rangle \cdots \langle \tilde{g}_{e_1,1} \rangle \cdots \langle \tilde{g}_{e_d,1} \rangle \cdots \langle \tilde{g}_{e_d,d} \rangle\]  

(152)

with the automorphism given by \((74)\) and \((75)\) and the evaluation function given by \((70)\) (where \(\tilde{ev}\) is replaced by \(ev\)) satisfying properties \((1)\) and \((2)\) of Lemma 9. Now consider the sub-difference ring \((L, \sigma)\) of \((\tilde{K}, \sigma)\) with \(L = \tilde{K}[\theta_{v_1,1}] \cdots [\theta_{v_d,1}] \cdots [\theta_{v_d,d}]\), which is a difference ring extension of \((\tilde{K}, \sigma)\), with the automorphism defined by

\[\sigma(\theta_{\ell,k}) = \gamma_{\ell,k} \theta_{\ell,k}\]  

(157)

for \(1 \leq \ell \leq d\) and \(1 \leq k \leq v_k\) where \(U = \langle \zeta \rangle\) is the multiplicative cyclic subgroup of \(\tilde{K}\) generated by a primitive \(\ell\)-th root of unity, \(\zeta \in \tilde{K}^\times\). Observe that the difference ring extension \((L, \sigma)\) of \((\tilde{K}, \sigma)\) with \((16)\) is a simple A-extension to which statement \((1)\) of Theorem 7 can be applied. Thus there is an R-extension \((K'[\theta], \sigma)\) of \((K', \sigma)\) with

\[\sigma(\theta) = \zeta' \theta\]  

(158)

of order \(\lambda'\) where \(K' = K'(\kappa_1, \ldots, \kappa_u)\), \(\zeta'\) is a primitive \(\lambda'\)-th root of unity in \(K'\) and \(K'\) is a finite algebraic field extension of \(\tilde{K}\). Note that the difference ring \((D, \sigma)\) where \(D\) is given by

\[D = K'[\tilde{y}_1,1] \cdots \langle \tilde{y}_{e_1,1} \rangle \cdots \langle \tilde{y}_{e_d,1} \rangle \cdots \langle \tilde{y}_{e_d,d} \rangle\]  

(159)

with the automorphism defined by \((75)\) is a \(\Pi\)-extension of \((K', \sigma)\). Thus by Lemma 3 it follows that the A-extension \((D[\theta], \sigma)\) of \((D, \sigma)\) with \((117)\) of order \(\lambda'\) is an R-extension. Note that the generators in the ring \(D[\theta]\) can be rearranged to get \((D, \sigma)\) where \(D = K'[\theta]\langle \tilde{y}_1,1 \rangle \cdots \langle \tilde{y}_{e_1,1} \rangle \cdots \langle \tilde{y}_{e_d,1} \rangle \cdots \langle \tilde{y}_{e_d,d} \rangle\) and \(\sigma\) is defined by \((117)\) and \((75)\). Since this rearrangement does not change the set of constants, \((D, \sigma)\) is an R-extension of \((K', \sigma)\). By statement \((1)\) of Proposition 3 \(D = e_0 D \oplus \cdots \oplus e_{\lambda'-1} D\) and by statement \((2)\) of the same proposition, \(e_k D = e_k D\) for \(0 \leq k < \lambda'\). Thus \(D = e_0 D \oplus e_1 D \oplus \cdots \oplus e_{\lambda'-1} D\) holds. Now we show that \(\phi : \tilde{G} \to D\) defined by \(\phi|_K = id_K\) with

\[\phi(\theta_{\ell,k}) = \tilde{y}_{\ell,k}\]  

(161)

\[\phi(\theta_{\ell,k}) = \beta_{\ell,k,0} e_0 + \cdots + \beta_{\ell,k,\lambda'-1} e_{\lambda'-1}\]  

(162)

where \(\beta_{\ell,k,i} = ev(\theta_{\ell,k}, \lambda' - 1 - i)\) for \(0 \leq i < \lambda'\) is a difference ring homomorphism. By statement \((1)\) of Theorem 7 \(\phi|_K\) which is defined by \((120)\) is a difference ring homomorphism. Since \(\phi\) maps \(\tilde{y}_{\ell,k}\) to itself, also \(\phi\) is a difference ring homomorphism. Furthermore, for all \(f \in \tilde{G}\) and all \(n \in \mathbb{Z}_{\geq 0}\), we have

\[ev'(f, n) = ev(\phi(f), n)\]  

(121)

Putting everything together, the map \(\varphi : A_d \to D\) with \(\varphi = \phi \circ \rho_d\) is a difference ring homomorphism. It is uniquely determined by \(\varphi|_K = id_K\) and

\[\varphi(y_{d},d) = \phi(\rho_d(y_{d},d)) = \gamma_{d,d} \tilde{y}_{1,1}^{v_{1,1,d}} \cdots \tilde{y}_{e_d,1}^{v_{e_d,1,d}} \]
with \( \gamma_{\ell,d} = \beta_{\ell,d,0} e_0 + \cdots + \beta_{\ell,d,N-1} e_{N-1} \in \mathbb{K}'[\theta] \). Furthermore, by (114) and (121) it follows that for all \( f \in A_d \) and all \( n \in \mathbb{Z}_{\geq 0} \) we get
\[
ev(f, n) = \ev'(\rho_d(f), n) = \ev(\phi(\rho_d(f)), n) = \ev(\varphi(f), n).
\]

Finally if \( K \) is strongly \( \sigma \)-computable, then by Lemma 9 the difference ring \((\widehat{\mathbb{A}}, \sigma)\) with (115) together with automorphism (74) and (75), evaluation function (76) (where \( \hat{\ev} \) is replaced by \( \ev \)) and the difference ring homomorphism \( \rho_d : A_d \to \widehat{\mathbb{A}} \) with (77) can be computed. Further, by Theorem 7 the difference ring \((\widehat{\mathbb{D}}[\theta], \sigma)\) with the automorphism \( \sigma(\theta) = \zeta' \theta \) and (75), the evaluation function (111) and the difference ring homomorphism \( \varphi : \widehat{\mathbb{A}} \to \widehat{\mathbb{D}} \) given by \((119)\) and \((120)\) can be computed. In particular, \( \varphi \) and all the components stated in the theorem can be given explicitly. \( \square \)

Example 13 (Cont. Example 9) (11). Take the AII-extension \((\mathbb{G}, \sigma)\) of \((\mathbb{K}, \sigma)\) with (85) constructed in Example 9 with the automorphism defined in (47) and (86), and consider the sub-difference ring \((\mathbb{K}[\bar{v}_{1,1}] | \bar{v}_{1,2}, \sigma)\) of \((\mathbb{G}, \sigma)\) with the automorphism \( \sigma \) given in (47), which is a simple A-extension of \((\mathbb{K}, \sigma)\) where \( \mathbb{K} = \mathbb{Q}(\sqrt{3}) \). Now we refine the construction from Example 9 by utilizing Example 11 Namely, we take the R-extension \((\mathbb{K} | \theta], \sigma)\) of \((\mathbb{K}, \sigma)\) of order 4 with the automorphism (104) and the evaluation function \( \hat{\ev} : \mathbb{K}[\theta] \times \mathbb{Z}_{\geq 0} \to \mathbb{K} \) given by (105) where \( \mathbb{K} = \mathbb{Q}(1, \sqrt{3}) \). Furthermore, take \((\mathbb{D}, \sigma)\) where \( \mathbb{D} = \mathbb{K}[\theta](y_{1,1})(y_{1,2})(y_{2,2})(y_{3,2}) \) with the automorphism and evaluation function given by (104) and (105) for the \( R \)-monomial \( \theta \) and (86) for the II-monomials \( y_{1,1,2,3} \). By Theorem 8 \((\mathbb{D}, \sigma)\) is an AII-extension of \((\mathbb{K}, \sigma)\) where the ring \( \mathbb{D} \) can be written as the direct sum \( \mathbb{D} = e_0 \mathbb{D} \oplus e_1 \mathbb{D} \oplus e_2 \mathbb{D} \oplus e_3 \mathbb{D} \) with \( \mathbb{D} = \mathbb{K}(y_{1,1})(y_{1,2})(y_{3,1})(y_{1,2})(y_{3,2}) \); here the idempotent elements \( e_k \) for \( 0 \leq k \leq 3 \) are defined by (106). Furthermore, the ring homomorphism \( \phi : \mathbb{G} \to \mathbb{D} \) defined by \( \phi|_{\mathbb{G}} = \text{id}_{\mathbb{G}} \) and (107) is a difference ring homomorphism.

Finally, consider the AP-extension \((\mathbb{A}', \sigma)\) of \((\mathbb{K}, \sigma)\) as given in Example 9 and consider the difference ring homomorphism \( \rho : \mathbb{A}' \to \mathbb{G} \) given in (57). Then with the difference ring homomorphism \( \varphi : \mathbb{A}' \to \mathbb{D} \) defined by \( \rho(\varphi(f)) \) for \( f \in \mathbb{A}' \) we get (113) for all \( n \in \mathbb{Z}_{\geq 0} \) and \( f \in \mathbb{A}' \). Given this explicit construction we can choose for instance \( g \in \mathbb{A}' \) defined in (57) that models \( \tilde{G}(n) \) given in (37). This means that \( \ev(g, n) = \tilde{G}(n) \) for all \( n \geq 0 \). Thus
\[
\tilde{g} := \varphi(g) = \frac{\varphi(y_{1,1}) \varphi(y_{3,1}) \varphi(y_{1,2}) \varphi(y_{2,2})}{\varphi(y_{1,1}) \varphi(y_{1,2}) \varphi(y_{4,2})} = \frac{(1 - i)(1 + i)^2 + 1}{2} \frac{y_{1,1} y_{2,1} y_{3,2}}{y_{2,1} y_{3,2}} \in \mathbb{D}
\]
yields for \( n \geq 0 \) the identity
\[
\tilde{G}(n) = \ev(g, n) = \ev(\tilde{g}, n) = \frac{1}{2} \frac{(1 - i)(1 + i)^2 + 1}{2} \frac{(\sqrt{3})^n (5)^n 2(\frac{n + 1}{2})}{2n 5(\frac{n - 1}{2})}.
\]

6.2. The solution for nested hypergeometric products

So far we have treated hypergeometric products over monic irreducible polynomials of finite nesting depth, say \( b \), that are \( \delta \)-defined for some \( \delta \in \mathbb{Z}_{\geq 0} \); see Definition 5. Given such hypergeometric products, it follows by Corollary 2 that we can construct an ordered multiple chain II-extension \((\overline{\mathbb{E}}_b, \sigma)\) of \((\mathbb{K}(x), \sigma)\) with \( \mathbb{K} = \mathbb{K}(\kappa_1, \ldots, \kappa_n) \) and
\[
\overline{\mathbb{E}}_b = \mathbb{K}(x)(\bar{z}_1) \cdots (\bar{z}_b) = \mathbb{K}(x)(\bar{z}_{p_1,1}) \cdots (\bar{z}_{p_1,b}) \cdots (\bar{z}_{p_s,1}) \cdots (\bar{z}_{p_s,b}).
\]
In particular, \((\overline{\mathbb{E}}_b, \sigma)\) is composed by the single chain II-extensions \((\overline{\mathbb{F}}_\ell, \sigma)\) of \((\mathbb{K}(x), \sigma)\) for \( 1 \leq \ell \leq p_1 \) with
\[
\overline{\mathbb{F}}_\ell = \mathbb{K}(x)(\bar{z}_{\ell,1}) \cdots (\bar{z}_{\ell,k}) \cdots (\bar{z}_{\ell,s_{\ell}}), \quad 1 \leq k \leq s_{\ell}
\]
given by the automorphism \( \sigma : \overline{\mathbb{F}}_\ell \to \overline{\mathbb{F}}_\ell \) defined by
\[
\sigma(\bar{z}_{\ell,k}) = \tilde{\alpha}_{\ell,k} \tilde{z}_{\ell,k} \quad \text{where} \quad \tilde{\alpha}_{\ell,k} = f \tilde{z}_{\ell,1} \cdots \tilde{z}_{\ell,k-1} \in (\mathbb{K}(x)^*)^{\mathbb{K}(x)(\bar{z}_{\ell,1}) \cdots (\bar{z}_{\ell,k-1})}
\]
and the evaluation function $\tilde{ev} : \tilde{F}_\ell \times \mathbb{Z}_{\geq 0} \to \tilde{K}$ defined by

$$\tilde{ev}(\tilde{z}_{\ell,k}, n) = \prod_{j=\delta}^{n} \tilde{ev}(\tilde{a}_{\ell,k}, j-1).$$

(125)

On the other hand, geometric products over the contents were treated in Subsection 5.2. In Theorem 8 we constructed a simple $R\Pi$-extension $(\mathbb{D}, \sigma)$ of $(\tilde{K}, \sigma)$ with $\tilde{K} = \tilde{K}(\kappa_1, \ldots, \kappa_u)$ where $\tilde{K}$ is a finite algebraic field extension of $K$ in which the geometric products can be modelled. To accomplish this task, we set up a ring of the form

$$\mathbb{D} = \tilde{K}[\vartheta](\tilde{y}_{1,1}) \ldots (\tilde{y}_{1,d}) \ldots (\tilde{y}_{e_d,d})$$

(126)

with

(a) the automorphism $\sigma : \mathbb{D} \to \mathbb{D}$ defined by

$$\sigma(\vartheta) = \zeta \vartheta,$$

(127)

$$\sigma(\tilde{y}_{\ell,k}) = \tilde{\gamma}_{\ell,k} \tilde{y}_{\ell,k}$$

(128)

where $\zeta \in \tilde{K}^*$ is a $\lambda$-th root of unity and $\tilde{\gamma}_{\ell,k} = \tilde{h}_\ell \tilde{y}_{\ell,1} \cdots \tilde{y}_{\ell,k-1} \in (\tilde{K}^*)_{\tilde{K}}(\tilde{y}_{1,1}) \ldots (\tilde{y}_{1,d})$ for $1 \leq k \leq d$ and $1 \leq \ell \leq e_k$ and

(b) the evaluation function $\tilde{ev} : \mathbb{D} \times \mathbb{Z}_{\geq 0} \to \tilde{K}$ defined by

$$\tilde{ev}(\vartheta, n) = \prod_{j=1}^{n} \zeta,$$

(129)

$$\tilde{ev}(\tilde{y}_{\ell,k}, n) = \prod_{j=1}^{n} \tilde{ev}(\tilde{\gamma}_{\ell,k}, j-1).$$

(130)

In particular, by reordering we obtain the difference ring extension $(\tilde{A}_d, \sigma)$ of $(\tilde{K}, \sigma)$ with

$$\tilde{A}_d = \tilde{K}(\tilde{y}_1) \cdots (\tilde{y}_d) = \tilde{K}(\tilde{y}_{1,1}) \ldots (\tilde{y}_{e_1,1}) \ldots (\tilde{y}_{1,d}) \ldots (\tilde{y}_{e_d,d}),$$

(131)

the automorphism (128) and the evaluation function (130) which is a sub-difference ring of $(\mathbb{D}, \sigma)$. Furthermore, it is an ordered multiple chain $\Pi$-extension of $(\tilde{K}, \sigma)$ and is composed by the single chain $\Pi$-extensions $(\tilde{K}_\ell, \sigma)$ of $(\tilde{K}, \sigma)$ where $\tilde{K}_\ell = \tilde{K}(\tilde{y}_{\ell,1}) \cdots (\tilde{y}_{\ell,e_\ell})$ for $1 \leq \ell \leq e_1$.

Putting the two difference rings $(\tilde{E}_b, \sigma)$ with (123) and $(\mathbb{D}, \sigma)$ with (126) together, we will obtain a difference ring in which we can model any finite set of hypergeometric product expressions of finite nesting depth coming from $\text{Prod}_{\mathbb{E}}(\mathbb{K}(x))$. Before we can complete this final argument, we have to take care that the two combined extensions yield again an $R\Pi$-extension. Here we utilize the following result from [OS18] Lemma 5.6 that holds for single nested $R\Pi$-extensions.

**Lemma 12.** Let $(\mathbb{K}(x), \sigma)$ be the rational difference field with $\sigma(x) = x + 1$ and let $(\mathbb{K}(x)(z_1) \ldots (z_s), \sigma)$ be a $\Pi$-extension of $(\mathbb{K}(x), \sigma)$ with $\frac{\sigma(z_k)}{z_k} \in \mathbb{K}[x] \setminus \mathbb{K}$. Further, let $\mathbb{K}'$ be an algebraic field extension of $\mathbb{K}$ and let $(\mathbb{K}'(y_1) \ldots (y_w), \sigma)$ be a $\Pi$-extension of $(\mathbb{K}', \sigma)$ with $\frac{\sigma(y_1)}{y_1} \in \mathbb{K}' \setminus \{0\}$. Then the difference ring $(\mathbb{E}, \sigma)$ with $\mathbb{E} = \mathbb{K}'(x)(y_1) \cdots (y_w)\langle z_1 \rangle \cdots \langle z_s \rangle$ is a $\Pi$-extension of $(\mathbb{K}'(x), \sigma)$. Furthermore, the $A$-extension $(\mathbb{E}[\vartheta], \sigma)$ of $(\mathbb{E}, \sigma)$ with $\sigma(\vartheta) = \zeta \vartheta$ of order $\lambda$ is an $R$-extension.

Namely, we can enhance the above lemma to nested $R\Pi$-extensions.
Corollary 3. Let \((\mathbb{K}(x), \sigma)\) be a rational difference field over \(\mathbb{K}\) with \(\sigma(x) = x + 1\) and let the difference ring \((\mathbb{H}_0, \sigma)\) with \(\mathbb{H}_0 = \mathbb{K}(x)_Z\) be an ordered multiple chain II-extension of \((\mathbb{K}(x), \sigma)\) with the automorphism \(\mathfrak{H}\) \(\mathfrak{H}\) \(\mathfrak{H}\) of \(
abla\). Further, let \(\mathbb{K}\) be an algebraic field extension of \(\overline{\mathbb{K}}\) and let the difference ring \((\mathbb{H}_d, \sigma)\) with \(\mathbb{H}_d = \mathbb{K}(x)_Z\) be the ordered multiple chain II-extension of \((\mathbb{K}, \sigma)\) with the automorphism \(\mathfrak{H}\). Then the difference ring \((\mathbb{E}, \sigma)\) with
\[
\mathbb{E} = \mathbb{K}(x)_Z \langle \tilde{y}_1, \ldots, \tilde{y}_d \rangle \langle \tilde{z}_1, \ldots, \tilde{z}_b \rangle
\]
where \(\langle \tilde{y}_i \rangle = \langle \tilde{y}_{i1}, \ldots, \tilde{y}_{i e_i} \rangle \) for \(1 \leq i \leq d\) and \(\langle \tilde{z}_k \rangle = \langle \tilde{z}_{k1}, \ldots, \tilde{z}_{k e_k} \rangle \) for \(1 \leq k \leq b\) is an ordered multiple chain II-extension of \((\mathbb{K}(x), \sigma)\). Furthermore, the \(\Lambda\)-extension \((\mathbb{E}, \sigma)\) of \((\mathbb{E}, \sigma)\) where \(\mathbb{E} = \mathbb{E}[\sigma]\) with \(\mathfrak{H}\) of order \(\lambda\) is an \(R\)-extension.

Proof. Take the II-extensions \((\mathbb{H}_1, \sigma)\) of \((\mathbb{K}(x), \sigma)\) with \(\mathbb{H}_1 = \mathbb{K}(x)_Z \langle \tilde{z}_1, \ldots, \tilde{z}_b \rangle\) and \((\mathbb{A}_1, \sigma)\) of \((\mathbb{K}, \sigma)\) with \(\mathbb{A}_1 = \mathbb{K}(x)_Z \langle \tilde{y}_1, \ldots, \tilde{y}_d \rangle\) which are both of monomial depth 1. By Lemma 12 the difference ring \((\mathbb{E}_1, \sigma)\) with \(\mathbb{E}_1 = \mathbb{K}(x)_Z \langle \tilde{y}_1, \ldots, \tilde{y}_d \rangle \langle \tilde{z}_1, \ldots, \tilde{z}_b \rangle\) is a II-extension of \((\mathbb{K}(x), \sigma)\) of monomial depth 1. Consider the ordered multiple chain P-extension \((\mathbb{E}, \sigma)\) of \((\mathbb{E}(x), \sigma)\) with \(\mathbb{E}\) which is composed by the single chain II-extensions in the ordered multiple chains \((\mathbb{H}_d, \sigma)\) and \((\mathbb{A}_d, \sigma)\). By Theorem 3 together with Lemma 3 it follows that \((\mathbb{E}, \sigma)\) is a II-extension of \((\mathbb{K}(x), \sigma)\). The quotient field of \(\mathbb{E}\) gives the rational function field \(\mathbb{H} = \mathbb{K}(x)_Z \langle \tilde{y}_1, \ldots, \tilde{y}_d \rangle \langle \tilde{z}_1, \ldots, \tilde{z}_b \rangle\) and one can extend the automorphism \(\sigma\) from \(\mathbb{E}\) to \(\mathbb{H}\) accordingly. Then by Lemma 3 \((\mathbb{H}, \sigma)\) is a \(\Pi\Sigma\)-field over \(\mathbb{K}\). Thus by Lemma 3 the \(\Lambda\)-extension \((\mathbb{H}[\sigma], \sigma)\) of \((\mathbb{H}, \sigma)\) of order \(\lambda\) with the automorphism \(\mathfrak{H}\) is an \(R\)-extension. Therefore \(\text{const}(\mathbb{H}[\sigma]) = \mathbb{K}\) with \(\mathbb{K} \subseteq \mathbb{E}[\sigma] \subseteq \mathbb{H}\) it follows that \(\text{const}(\mathbb{E}[\sigma]) = \mathbb{K}\). But this implies that the \(\Lambda\)-extension \((\mathbb{E}[\sigma], \sigma)\) of \((\mathbb{E}, \sigma)\) of order \(\lambda\) with the automorphism \(\mathfrak{H}\) is an \(R\)-extension. \(\square\)

Finally, we arrive at the following main result.

Theorem 9. Let \(\mathbb{K} = \mathbb{K}(\kappa_1, \ldots, \kappa_n)\) be a rational function field, and let \((\mathbb{K}(x), \sigma)\) with \(\sigma(x) = x + 1\) be a rational difference field with the evaluation function \(\text{ev} : \mathbb{K}(x) \times \mathbb{Z}_{\geq 0} \to \mathbb{K}\) defined by \((\mathbb{H})\), and the \(Z\)-function given by \((\mathbb{H})\). Suppose we are given a finite set of hypergeometric product expressions
\[
\{A_1(n), \ldots, A_m(n)\} \subseteq \text{Prod}_{\mathbb{H}}(\mathbb{K}(x))
\]
of nesting depth at most \(d\) for some \(d \in \mathbb{Z}_{\geq 0}\). Then there is a \(\delta \in \mathbb{Z}_{\geq 0}\) and an \(\Pi\Sigma\)-extension \((\mathbb{E}, \sigma)\) of \((\mathbb{K}(x), \sigma)\) of monomial depth at most \(d\) where \(\mathbb{E}\) is a finite algebraic field extension of \(\mathbb{K}\) equipped with an evaluation function \(\text{ev} : \mathbb{E} \times \mathbb{Z}_{\geq 0} \to \mathbb{K}\) with respect to \(\delta\) with the following properties:

1. The map \(\tau : \mathbb{E} \to \mathbb{S}(\mathbb{K})\) with \(\tau(f) = (\text{ev}(f, n))_{n \geq 0}\) is a \(\mathbb{K}\)-embedding.

2. There are elements \(a_1, \ldots, a_e \in \mathbb{E}^*\) such that for \(j\) with \(1 \leq j \leq e\) and for all \(n \geq \delta\) we have
\[
A_j(n) = \text{ev}(a_j, n).
\]

If \(K\) is a strongly \(\sigma\)-computable, such a \(\delta\), \((\mathbb{E}, \sigma)\) with the evaluation function \(\text{ev}\), and the \(a_1, \ldots, a_m \in \mathbb{E}\) can be computed.

Proof.

(a) We are given the hypergeometric product expressions in \((\mathbb{H})\) where
\[
\sum_{\nu = (\nu_1, \ldots, \nu_e) \in S_j} a_{\nu}(n) P_1(n)^{\nu_1} \cdots P_e(n)^{\nu_e}
\]
with \(S_j \subseteq \mathbb{Z}_{\geq 0}\) finite, \(a_{\nu}(n) \in \mathbb{K}(x)\) and \(P_1(n), \ldots, P_e(n) \in \text{Prod}_{\mathbb{H}}(\mathbb{K}(x))\). Now we follow the construction in Proposition 2. There we can take a \(\delta \in \mathbb{Z}_{\geq 0}\) and construct for all \(1 \leq j \leq e\), \(c_j \in \mathbb{K}^*\), rational functions \(r_j \in \mathbb{K}(x)^*\), 1-refined geometric product expressions \(\tilde{G}_j(n) \in \text{Prod}_{\mathbb{H}}(\mathbb{K})\)
and $\delta$-refined hypergeometric product expressions in shift-coprime product representation form $\tilde{H}_j(n) \in \text{ProdM}_n(\mathbb{K}(x))$ such that

$$P_j(n) = \tilde{c}_j \tilde{r}_j(n) \tilde{G}_j(n) \tilde{H}_j(n) \neq 0$$

holds for all $n \geq \max(0, \delta - 1)$.

(b) For the hypergeometric product expressions $\tilde{H}_1(n), \ldots, \tilde{H}_c(n)$ in (134) we have

$$\tilde{H}_i(n) = \tilde{H}_{i,1}(n)^{n_{i,1}} \cdots \tilde{H}_{i,l_i}(n)^{n_{i,l_i}}$$

for some $l_i \in \mathbb{Z}_{\geq 0}$ with $n_{i,j} \in \mathbb{Z}$ for $1 \leq j \leq l_i$ where all the arising hypergeometric products $\tilde{H}_{i,j}(n)$ are $\delta$-refined and in shift-coprime product representation form. By Corollary 2 we can construct an ordered multiple chain II-extension ($\tilde{H}_i, \sigma$) of $(\mathbb{K}(x), \sigma)$ with (123) which is composed by the single chain II-extensions $(\tilde{F}_\ell, \sigma)$ of $(\mathbb{K}(x), \sigma)$ with $1 \leq \ell \leq p_1$ for some $p_1 \in \mathbb{Z}_{\geq 0}$ with $\tilde{F}_\ell = \mathbb{K}(x)\langle \tilde{z}_{\ell,1}, \tilde{z}_{\ell,2}, \ldots, \tilde{z}_{\ell,s_\ell} \rangle$, the automorphism $\sigma : \tilde{F}_\ell \rightarrow \tilde{F}_\ell$ given in (124) and the evaluation function $\tilde{ev} : \tilde{F}_\ell \times \mathbb{Z}_{\geq 0} \rightarrow \mathbb{K}$ defined by $\tilde{ev}|_{\mathbb{K}(x) \times \mathbb{Z}_{\geq 0}} = \text{ev}$ and (125). In particular, there are $\nu_{i,j}, \mu_{i,j}$ such that $\tilde{H}_{i,j}(n) = \text{ev}(\tilde{z}_{\nu_{i,j}, \mu_{i,j}}, n)$ holds for all $n \geq \max(0, \delta - 1)$. Thus we can take $\tilde{h}_i = \tilde{z}_{\nu_{i,1}, \mu_{i,1}} \cdots \tilde{z}_{\nu_{i,l_i}, \mu_{i,l_i}} \in \tilde{H}_i$ with

$$\tilde{ev}((\tilde{h}_i, n)) = \tilde{H}_j(n) \quad \forall n \geq \delta.$$  

(c) Next we treat the geometric product expressions $\tilde{G}_1(n), \ldots, \tilde{G}_c(n)$ in (134). Following Remark 4 we can construct a multiple chain P-extension ($\tilde{A}, \sigma$) of $(\mathbb{K}, \sigma)$ where the bases are from $\mathbb{K}^*$ such that there are $g_1, \ldots, g_c \in \tilde{A}$ with $\text{ev}(g_i, n) = G_i(n)$ for all $n \in \mathbb{Z}_{\geq 0}$. Then by Theorem 5 we can construct an RII-extension $(\tilde{A}, \sigma)$ of $(\mathbb{K}, \sigma)$ with (126) together with the automorphism $\sigma : \tilde{A} \rightarrow \tilde{A}$ given in (127) and (128), with the evaluation function $\tilde{ev} : \tilde{A} \times \mathbb{Z}_{\geq 0} \rightarrow \mathbb{K}$ defined by $\tilde{ev}(c, n) = c$ for all $c \in K$, $n \in \mathbb{Z}_{\geq 0}$, (129) and (130), and with a difference ring homomorphism $\varphi : \tilde{A}_d \rightarrow \tilde{A}$ such that for all $f \in \tilde{A}$ and $n \in \mathbb{Z}_{\geq 0}$ we have (113). Thus for $\tilde{g}_j := \varphi(g_j)$ with $1 \leq j \leq e$ we get

$$\tilde{ev}(\tilde{g}_j, n) = \tilde{ev}(\varphi(g_j), n)$$

and (134)

$$\text{ev}(g_j, n) = \tilde{G}_j(n) \quad \forall n \geq 0.$$  

(d) By Corollary 3 we can merge these two difference rings to obtain an RII-extension $(\tilde{E}, \sigma)$ of $(\tilde{\mathbb{K}}(x), \sigma)$ with (132) and the automorphism $\sigma : \tilde{E} \rightarrow \tilde{E}$ and the evaluation function $\tilde{ev} : \tilde{E} \times \mathbb{Z}_{\geq 0} \rightarrow \mathbb{K}$ defined accordingly. As $(\tilde{E}, \sigma)$ is an RII-extension of the rational difference field $(\tilde{\mathbb{K}}(x), \sigma)$, it follows by Theorem 2 that $\tilde{r} : \tilde{E} \rightarrow \tilde{S}(\mathbb{K})$ defined by (67) is a $\mathbb{K}$-embedding. For $1 \leq j \leq e$, define

$$p_j := \tilde{c}_j \tilde{r}_j \tilde{g}_j \tilde{h}_j \in \tilde{E}.$$  

With $\tilde{ev}(\tilde{r}_j, n) = \tilde{r}_j(n)$ and the evaluations of $\tilde{h}_j$ and $\tilde{g}_j$ given in (135) and (136) together with (134) it follows that for all $1 \leq j \leq e$ and for all $n \geq \max(0, \delta - 1)$ we have

$$P_j(n) = \tilde{c}_j \tilde{r}_j(n) \tilde{G}_j(n) \tilde{H}_j(n) = \tilde{ev}(\tilde{c}_j, n)\tilde{ev}(\tilde{r}_j, n)\tilde{ev}(\tilde{g}_j, n)\tilde{ev}(\tilde{h}_j, n) = \tilde{ev}(\tilde{c}_j \tilde{r}_j \tilde{g}_j \tilde{h}_j, n) = \tilde{ev}(p_j, n).$$

(e) Finally, we can define $a_j = \sum_{\nu_1, \ldots, \nu_e} p_1^{\nu_1} \cdots p_e(n)^{\nu_e} \in \tilde{E}$ for $1 \leq j \leq m$ and get $\text{ev}(a_j, n) = A_j(n)$ for all $n \geq \max(\delta - 1, 0)$. Observe that if $\mathbb{K}$ is strongly $\sigma$-computable, all the ingredients delivered by Corollary 2 and Theorem 8 can be computed. In particular, $(\tilde{E}, \sigma)$, and $\sigma$ with $\tilde{ev}$ and $a_1, \ldots, a_m$ can be computed explicitly.

As a consequence we are now in the position to solve Problem RPE as follows.
Corollary 4. Let \( A(n) \in \text{ProdE}_n(\mathbb{K}(x)) \) with \( x \). For \( A_1(n) = A(n) \) with \( m = 1 \) let \( \delta \in \mathbb{Z}_{\geq 0} \), \((\mathbb{E}, \sigma)\) with the evaluation function \( \text{ev} \) and the \( a := a_1 \in \mathbb{E} \) be the ingredients as provided in Theorem 9. In particular, let \( \mathbb{E} = \mathbb{K}(x)[\vartheta](p_1 \ldots p_s) \) where \( \vartheta \) is the \( \mathbb{R} \)-monomial with \( \sigma(\vartheta) = \varsigma \vartheta \) and let \( p_1, \ldots, p_s \) be the \( \Pi \)-monomials. Furthermore, let \( a = \sum_{v=(\nu_1, \ldots, \nu_s) \in \hat{S}} b_v(n) \vartheta^{\nu} p_1^{\mu_1} \cdots p_s^{\mu_s} \) with \( \hat{S} \subseteq \{0, \ldots, \lambda - 1\} \times \mathbb{N}^s \) finite and \( a_v(x) \in \mathbb{K}(x) \) for \( v \in \hat{S} \). Then the following holds:

1. \( \text{ev}(\vartheta, n) = \xi^n \) for all \( n \in \mathbb{Z}_{\geq 0} \); furthermore, for \( 1 \leq i \leq s \) we have \( \text{ev}(p_i, n) = Q_i(n) \) for all \( n \geq 0 \) where \( Q_i(n) \in \text{Prod}_n(\mathbb{K}(x)) \).

2. For \( B(n) = \sum_{v=(\nu_1, \ldots, \nu_s) \in \hat{S}} b_v(n) (\xi^n)^{\mu_1} Q_1(n)^{\mu_1} \cdots Q_s(n)^{\mu_s} \) we have \( A(n) = B(n) \) for all \( n \geq \delta \).

3. The subring

\[
\tau(\mathbb{K}(x))(\xi^n | (Q_1(n))_{n \geq 0}, (Q_1(n)^{-1})_{n \geq 0}, \ldots, (Q_s(n))_{n \geq 0}, (Q_s(n)^{-1})_{n \geq 0})
\]

of \( \mathbb{S}(\mathbb{K}) \) forms a Laurent polynomial ring extension of \( \tau(\mathbb{K}(x))(\xi^n |_{n \geq 0}) \). Thus the sequences produced by \( Q_1(n), \ldots, Q_s(n) \) are algebraically independent among each other over \( \tau(\mathbb{K}(x))(\xi^n |_{n \geq 0}) \).

4. We have that \( A(n) = 0 \) for all \( n \geq d \) for some \( d \in \mathbb{Z}_{\geq 0} \) if and only if \( a = 0 \) if and only \( B(n) \) is the zero-expression. If this holds, \( A(n) = 0 \) for all \( n \geq \delta \).

Proof.

1. Note that \((\mathbb{E}, \sigma)\) is a multiple chain \( \Pi \)-extension of \((\mathbb{K}(x)[\vartheta], \sigma)\) over \( \mathbb{K}(x) \) equipped by an evaluation function given by the iterative application of Lemma 2. Thus statement (1) follows.

2. By statement (2) of Theorem 9 it follows that \( \text{ev}(a, n) = A(n) \) holds for all \( n \geq \delta \). By definition of the \( \text{ev} \) function we have \( \text{ev}(a, n) = B(n) \). Thus statement (2) holds.

3. Since \( \tau : \mathbb{E} \to \mathbb{S}(\mathbb{K}) \) with \( \tau(f) = (\text{ev}(\varphi, f(n))_{n \geq 0} \) is an injective difference ring homomorphism by statement (2) of Theorem 9, statement (3) follows.

4. Since \( \tau \) is injective, it follows that

\[
0 = A(n) \text{ for all } n \geq d \text{ for some } d \in \mathbb{Z}_{\geq 0} \quad \text{item (2)} \quad 0 = B(n) = \text{ev}(a, n) \text{ for all } n \geq \max(d, \delta) \\
\tau \text{ injective} \quad a = 0 \quad B(n) \text{ is the zero-expression.}
\]

If this is the case, \( A(n) = B(n) = 0 \) for all \( n \geq \delta \). This proves the last statement.

\[\square\]

Example 14 (Cont. Examples 1, 5, 7, 8, 9). Let \((\mathbb{K}(x), \sigma)\) be the rational difference field with \( \mathbb{K} = \mathbb{Q}(\sqrt{3}) \) equipped with the field automorphism \( \sigma : \mathbb{K}(x) \to \mathbb{K}(x) \) and the evaluation function \( \text{ev} : \mathbb{K}(x) \times \mathbb{Z}_{\geq 0} \to \mathbb{K} \) defined by \( \sigma(x) = x + 1 \) and \( 46 \) respectively. Given the nesting depth 2 hypergeometric product \( P(n) \) with \( 8 \) in Example 1, we computed the following 3-refined shift-coprime product representation form:

\[
P(n) = \hat{\varsigma} \tau(n) \hat{G}(n) \hat{H}(n) \tag{137}
\]

where \( \hat{\varsigma}, r(n), \hat{G}(n) \) and \( \hat{H}(n) \) are given in \( 35 \), \( 36 \), \( 37 \), and \( 38 \) respectively. In particular, \( 137 \) holds for all \( n \in \mathbb{Z}_{\geq 0} \) with \( n \geq \delta - 1 = 2 \).
In Example 8 we constructed the ordered multiple chain II-extension \((\hat{\mathbb{H}}, \sigma)\) of \((\mathbb{K}(x), \sigma)\) with \(\hat{\mathbb{H}} = \mathbb{K}(x)(z_{1,1}, z_{2,2}(z_{1,2})\) which is composed by the single chain II-extensions of \((\mathbb{K}(x), \sigma)\) defined in items [7] and [8]. The automorphism and the evaluation function were defined as given in [54] and [55]. In particular, the hypergeometric product expression \(\hat{H}(n)\) is modelled by the expression \(\hat{h} = z_{1,1}^h z_{2,1} z_{1,2}\) where \(\hat{h} = h\) is taken from [57]. Furthermore, we constructed the RII-extension \((\mathbb{D}, \sigma)\) of \((\hat{\mathbb{K}}, \hat{\sigma})\) with \(\hat{\mathbb{K}} = \mathbb{K}(\sqrt[3]{3}, i)\) equipped with the evaluation function \(\hat{\nu}\) from Example 13. There \(\hat{\mathbb{E}} = \hat{\mathbb{E}}(\hat{\nu})\) is isomorphic to \((\mathbb{K}, \mathbb{S})\). In particular, the hypergeometric product expression \(\hat{\nu}(\hat{e})\) is composed by the single chain Π-extensions of \((\mathbb{K}, \mathbb{S})\) where \(\hat{h} = h\) is is taken from (57). The Mathematica package — NestedProduct — is modelled by \((\mathbb{D}, \sigma)\) and \((\mathbb{E}, \sigma)\), we get the RII-extension \((\mathbb{E}, \sigma)\) of \((\hat{\mathbb{K}}, \hat{\sigma})\) with \(\mathbb{E} = \hat{\mathbb{K}}(\mathbb{D})(\mathbb{G})\langle y_{1,1}, y_{2,1}, y_{1,2}, y_{2,2}(\hat{y}_{1,2}),(\hat{y}_{1,1})\rangle\langle z_{1,1}, z_{2,1}, z_{1,2}\rangle\) where the automorphism \(\mathbb{E} : \mathbb{E} \to \mathbb{E}\) and the evaluation function \(\hat{\nu} : \mathbb{E} \times \mathbb{Z}_{\geq 0} \to \hat{\mathbb{K}}\) are defined by (104), (105), (86), (54), and (55). Following the proof of Theorem 9 we set \(p := \hat{c} \hat{r} \hat{y} \hat{h}\) and get for all \(n \geq 2\) the simplification

\[
P(n) = \hat{\nu}(p, n) = -\frac{254}{432} \hat{\nu}(r, n) \hat{\nu}(y, n) \hat{\nu}(h, n) \]

\[
= -\frac{254}{432} (n - 1)^3 n(n + 1)(n + 2) \frac{1}{2} (1 - i)^6 (i^5 + 1) (\sqrt{3})^n (5^n)^2 \left(\frac{n}{2}\right)^2 2 \left(\frac{n}{2}\right) 2^{n} 5^{n+1} \]

\[
\left(\prod_{k=3}^{n} (k - 2)\right)^3 \left(\prod_{k=3}^{n} (k + \frac{1}{2})\right) \left(\prod_{j=3}^{n} (j - 2)\right). \]

Based on this representation in an RII-extension, we can extract the following extra property. Since \(\tau : \mathbb{E} \to \mathbb{S}(\mathbb{K})\) is a \(\mathbb{K}\)-embedding, it follows that the sub-difference ring \((R, S)\) of \((\mathbb{S}(\mathbb{K}), S)\) with

\[
R = \tau(\mathbb{K}(x))[i^n]_{n \geq 0}, (Q_1(n))_{n \geq 0}, (Q_1(n)^{-1})_{n \geq 0} \ldots (Q_8(n))_{n \geq 0}, (Q_8(n)^{-1})_{n \geq 0}\]

and

\[
Q_1(n) = (\sqrt{3})^n, \quad Q_2(n) = 2^n, \quad Q_3(n) = 5^n, \quad Q_4(n) = \prod_{k=1}^{n} (k - 2), \]

\[
Q_5(n) = \prod_{k=3}^{n} (k + \frac{1}{2}) \quad Q_6(n) = 2^{\left(\frac{n+1}{2}\right)}, \quad Q_7(n) = 5^{\left(\frac{n+1}{2}\right)}, \quad Q_8(n) = \prod_{k=3}^{n} (j - 2) \]

is isomorphic to \((\mathbb{E}, \sigma)\). In particular, we can conclude that \(R\) is a Laurent polynomial ring extension of the ring \(G = \tau(\mathbb{K}(x))[i^n]_{n \geq 0}\). In a nutshell, the sequences generated by the products \(Q_1(n), \ldots, Q_8(n)\) are algebraically independent among each other over the ring \(G\).

6.3. The Mathematica package — NestedProducts

In the following we will demonstrate how our tools can be activated with the help of the Mathematica package NestedProduct. We start with the nested hypergeometric product expression

\[
A(n) = \frac{1}{2} \prod_{k=1}^{n-1} \frac{1}{36} \left(\prod_{i=1}^{k-1} \frac{(i + 1)(i + 2)}{4(2i + 3)^2}\right) \in \text{ProdE}(\mathbb{Q}(x)) \quad (138)
\]

from [Kau18] Example 3 which was guessed using the Mathematica package RATE written by Christian Krattenthaler; see [Kra97].
However, for more complicated expression such a representation can be rather challenging. Taking the inner product as the first $\Pi$-monomial and the outermost product as the second $\Pi$-monomial.

Note that one could have represented the product in (138) directly within a $\Pi$-extension by simply applying the command ProductReduce to $A$ we solve Problem RPE and get the result

$$\text{Out}[5] = 1$$

After loading the package

```mathematica
<<NestedProducts.m
```

```plaintext
NestedProducts — A package by Evans Doe Ocansey — © RISC
```

into Mathematica, we define the product with the command

```mathematica
A = 1/2 * FormalProduct[1/36 * FormalProduct[(1 + n) + (2/n), {i, 1, k - 1}], {k, 1, n - 1}];
```

Here FormalProduct[f, {k, a, n+b}] (as shortcut one can use FProduct) defines a nested product $\prod_{k=a}^{n+b} f$ where $a, b$ are integers and the multiplicand $f$, free of $n$, must be an expression in terms of nested products whose outermost upper bounds are given by $k$ or an integer shift of $k$. Then applying the command ProductReduce to $A$ we solve Problem RPE and get the result

```mathematica
A = FormalProduct[f, {i, 1, k}];
```

$$\text{Out}[3] = 9 (2^n) (\prod_{k=1}^{n} (k + \frac{3}{2}))^4 (\prod_{i=1}^{n} (i + 1)^2) \left(\prod_{k=1}^{n} (k + 1) \prod_{i=1}^{n} (i + \frac{3}{2})\right)^2 (2^n + 3^n, \prod_{k=1}^{n} (k + \frac{3}{2}), \prod_{k=1}^{n} (k + 1), 2^n = \prod_{k=1}^{n} (k + \frac{3}{2}), \prod_{i=1}^{n} (i + \frac{3}{2}), \prod_{k=1}^{n} (i + 1). \quad (140)

Note that one could have represented the product in (138) directly within a II-extension by simply taking the inner product as the first II-monomial and the outermost product as the second II-monomial. However, for more complicated expression such a representation can be rather challenging.

The full capability of our machinery can be illustrated by combining, e.g., the expression (138) with the following related product (where one of the inner products is slightly modified):

```mathematica
B = ProductReduce[A]
```

$$\text{Out}[3] = 9 (2^n) (\prod_{k=1}^{n} (k + \frac{3}{2}))^4 (\prod_{i=1}^{n} (i + 1)^2) \left(\prod_{k=1}^{n} (k + 1) \prod_{i=1}^{n} (i + \frac{3}{2})\right)^2 (2^n + 3^n, \prod_{k=1}^{n} (k + \frac{3}{2}), \prod_{k=1}^{n} (k + 1), 2^n = \prod_{k=1}^{n} (k + \frac{3}{2}), \prod_{i=1}^{n} (i + \frac{3}{2}), \prod_{k=1}^{n} (i + 1). \quad (140)$$
By solving Problem RPE the expression can be rephrased in an RII-extension with the R-monomial \( i^n \) and the \( \Pi \)-monomials given in (140). In short, together with \( i^n \) the expression can be reduced again in terms of the algebraic independent products (140).

Similar expressions as given in (138) arise during challenging evaluations of determinants; see, e.g., MRRJS3, Zei96, Kra01. We expect that the new tools elaborated in this article will prove beneficial in related but more complicated product expressions.

We conclude this section by combining our tools from above with the summation package \texttt{Sigma}.

After loading in

\[
\texttt{In[6]} = \texttt{\textless \textless Sigma.m}
\]

we insert the sum

\[
\texttt{In[7]} = \texttt{mySum = SigmaSum}[-1 + (1 + k)(2 + k)^2 \prod_{j=1}^{k}(1 + j)^2 \prod_{j=1}^{k} j \prod_{i=1}^{k} (1 + i)^2 \\
- \frac{4}{3} (1 + 2(1 + k)^2 (3 + k) \prod_{j=1}^{k} - j(2 + j)) \prod_{j=1}^{k} j \prod_{i=1}^{k} - i(2 + i), \{k, 1, n\}];
\]

Afterwards we can activate the available summation algorithms of \texttt{Sigma} with the function call \texttt{SigmaReduce} and succeed in eliminating the summation sign:

\[
\texttt{Out[8]} = 4 - \frac{1}{3} (1 + n)^5 (2 + n)^2 \left(-3 + (1 + i) \left(-i + (i^n)^2\right)(3 + n)i^n\right) (n!)^5 \left(\prod_{i=1}^{k} \prod_{j=1}^{k} j\right)^2.
\]

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In other words, we have derived the simplification

\[
\sum_{k=1}^{n} \left(-1 + (1 + k) (2 + k)^2 \prod_{j=1}^{k}(1 + j)^2 \prod_{j=1}^{k} j \prod_{i=1}^{k} (1 + i)^2 \\
- \frac{4}{3} (1 + 2(1 + k)^2 (3 + k) \prod_{j=1}^{k} - j(2 + j)) \prod_{j=1}^{k} j \prod_{i=1}^{k} - i(2 + i)\right) \\
= 4 - \frac{1}{3} (1 + n)^5 (2 + n)^2 \left(-3 + (1 + i) \left(-i + (i^n)^2\right)(3 + n)i^n\right) (n!)^5 \left(\prod_{i=1}^{k} \prod_{j=1}^{k} j\right)^2.
\]

We emphasize that the summand given in \texttt{ln[8]} has been transformed internally with the package \texttt{NestedProducts} to the form

\[
\frac{1}{3} (1 + k)^2 \left(\prod_{i=1}^{k} i\right)^3 \left(\prod_{i=1}^{k} \prod_{j=1}^{k} j\right)^2 \left(-3 + (-1 - i) (2 + k)i^k + (-1 + i) (2 + k) (i^k)^3 \\
+ (1 + k)^3 (2 + k)^2 \left(3 + (-1 + i) (3 + k)i^k + (-1 - i) (3 + k) (i^k)^3\right) \left(\prod_{i=1}^{k} j\right)^2\right).
\]

Then \texttt{Sigma} reads off the derived products and rephrases them directly to a tower of RII-extensions (without the exploitation of the available tools in \texttt{Sigma} that can check whether the constant field remains unchanged). Afterwards the underlying summation algorithms of \texttt{Sigma} are applied to derive the final result.
7. Conclusion

We enhanced non-trivially the ideas from [Sch05, Sch14, OS18] (related also to [AP10, CFFL11]) in order to solve Problem RPE in Theorem 9 and Corollary 4 above. There we cannot only reduce or simplify expressions in terms of hypergeometric products of nesting depth 1 but in terms of hypergeometric products of nesting depth $\geq 1$. More precisely, the expression can be reduced to an expression in terms of one root of unity product of the form $\zeta^n$ and hypergeometric products $Q_1(n), \ldots, Q_s(n) \in \text{Prod}_n(K(x))$ of arbitrary but finite nesting depth which are algebraically independent among each other. This latter property has been extracted from results elaborated in [Sch17] (which are inspired by [PS97]). Combined with the existing difference ring algorithms for symbolic summation [Kar81, Sch01, Sch16] this yields a complete summation machinery to reduce and simplify nested sums over hypergeometric products of arbitrary but finite nesting depth.

A natural future task is to enhance this combined toolbox of the packages NestedProducts and Sigma further and to tackle, e.g., definite summation problems. In particular, the interaction with the available creative telescoping algorithms [Sch07a, Sch08, Sch10b, Sch15] and recurrence solving algorithms [ABPS20] should be explored further.

Following the ideas from [OS18] one might extend the above machinery to the class of nested $q$-hypergeometric products covering also the multibasic and mixed case [BP99]. Another open task is to combine the above ideas with contributions from [Sch20] (based on Smith normal form calculations) to find optimal representations of such nested products. This means that in the output expression the order $\lambda$ of the primitive root of unity $\zeta$ in $\zeta^n$ and the number $s$ of algebraically independent products $Q_1(n), \ldots, Q_s(n)$ should be minimized.

Finally, it would be interesting to see if the class of hypergeometric products of finite nesting depth can be generalized further to products of the form $[\square]$ where in the multiplicands the products do not appear only in form of Laurent polynomial expressions.

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