Equivalences on Acyclic Orientations

Matthew Macauley$^{1,3}$, Henning S. Mortveit$^{2,3}$

$^1$ Department of Mathematics, UCSB
$^2$ Department of Mathematics, Virginia Tech
$^3$ NDSSL, VBI, Virginia Tech, e-mail: {macauley,henning.mortveit}@vt.edu

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Abstract  The cyclic and dihedral groups can be made to act on the set $\text{Acyc}(Y)$ of acyclic orientations of an undirected graph $Y$, and this gives rise to the equivalence relations $\sim_\kappa$ and $\sim_\delta$, respectively. These two actions and their corresponding equivalence classes are closely related to combinatorial problems arising in the context of Coxeter groups, sequential dynamical systems, the chip-firing game, and representations of quivers.

In this paper we construct the graphs $C(Y)$ and $D(Y)$ with vertex sets $\text{Acyc}(Y)$ and whose connected components encode the equivalence classes. The number of connected components of these graphs are denoted $\kappa(Y)$ and $\delta(Y)$, respectively. We characterize the structure of $C(Y)$ and $D(Y)$, show how $\delta(Y)$ can be derived from $\kappa(Y)$, and give enumeration results for $\kappa(Y)$. Moreover, we show how to associate a poset structure to each $\kappa$-equivalence class, and we characterize these posets. This allows us to create a bijection from $\text{Acyc}(Y)/\sim_\kappa$ to $\text{Acyc}(Y')/\sim_\kappa \cup \text{Acyc}(Y'')/\sim_\kappa$, where $Y'$ and $Y''$ denote edge deletion and edge contraction for a cycle-edge in $Y$, respectively, which in turn shows that $\kappa(Y)$ may be obtained by an evaluation of the Tutte polynomial at $(1,0)$.

1 Introduction

An acyclic orientation $O_Y$ of an undirected graph $Y$ induces a partial ordering on the vertex set $v[Y]$ by $i \leq_{O_Y} j$ if there is a directed path from $i$ to $j$
in \( O_Y \). A cyclic 1-shift (left) of a linear extension of \( O_Y \) corresponds to converting a source of \( O_Y \) into a sink. This source-to-sink operation on \( \text{Acyc}(Y) \) gives rise to the equivalence relation denoted \( \sim_\kappa \). Reversing a linear extension of \( O_Y \) corresponds to reflecting all edge orientations in \( O_Y \). The coarser equivalence relation on \( \text{Acyc}(Y) \) obtained through source-to-sink operations and reflections is denoted \( \sim_\delta \).

This paper is organized as follows. In Section 3 we construct the equivalence relations \( \sim_\kappa \) and \( \sim_\delta \), and graphs \( C(Y) \) and \( D(Y) \) that have vertex set \( \text{Acyc}(Y) \), and whose connected components correspond to \( \kappa \)- and \( \delta \)-equivalence classes, respectively. We let \( \kappa(Y) \) and \( \delta(Y) \) denote the number of equivalence classes. In Section 4 we study the structure and properties of the graphs \( C(Y) \) and \( D(Y) \), show how \( \delta(Y) \) can be determined from \( \kappa(Y) \), and give bounds for these quantities. In Section 5 we show that one may associate a poset to each \( \kappa \)-equivalence class and use this to establish the bijection

\[
\Theta : \text{Acyc}(Y)/\sim_\kappa \longrightarrow (\text{Acyc}(Y')/\sim_\kappa) \bigcup (\text{Acyc}(Y'')/\sim_\kappa),
\]

where \( Y' \) and \( Y'' \) are the graphs formed by deleting and contracting a non-bridge edge \( e \) of \( Y \), respectively. This leads to a new proof of a recursion relation for \( \kappa(Y) \) in [3]. Finally, in the summary, we discuss how these constructions arise in other areas of mathematics, such as sequential dynamical systems, Coxeter groups, the chip-firing game, and the representation theory of quivers.

2 Terminology and Background

Let \( Y \) be an undirected, simple and loop-free graph with vertex set \( v[Y] = \{1, 2, \ldots, n\} \) and edge set \( e[Y] \). We let \( S_Y \) denote the set of total orders (i.e., permutations) of \( v[Y] \). In [3] an equivalence relation \( \sim_Y \) is introduced on \( S_Y \) through the update graph \( U(Y) \) of \( Y \). The update graph has vertex set \( S_Y \), and two distinct vertices \( \pi = (\pi_i)_i \) and \( \pi' = (\pi'_i)_i \) are adjacent if they differ in exactly two consecutive elements \( \pi_k \) and \( \pi_{k+1} \) such that \( \{\pi_k, \pi_{k+1}\} \notin e[Y] \). The equivalence relation \( \sim_Y \) is defined by \( \pi \sim_Y \pi' \) if \( \pi \) and \( \pi' \) are connected in \( U(Y) \). We denote the equivalence class containing \( \pi \) by \( [\pi]_Y \), and set

\[
S_Y/\sim_Y = \{[\pi]_Y \mid \pi \in S_Y\}.
\]

This corresponds to partially commutative monoids as defined in [3], but restricted to fixed length permutations over \( v[Y] \) and with commutation relations encoded by non-adjacency in the graph \( Y \).

Orientations of \( Y \) are represented as maps \( O_Y : e[Y] \longrightarrow v[Y] \times v[Y] \), which may also be viewed as directed graphs. The set of acyclic orientations of \( Y \) is denoted \( \text{Acyc}(Y) \), and we set \( \alpha(Y) = |\text{Acyc}(Y)| \). Every acyclic orientation defines a partial ordering on \( v[Y] \) where the covering relations are \( i \leq_{O_Y} j \) if \( \{i, j\} \in e[Y] \) and \( O_Y(\{i, j\}) = (i, j) \). The set of linear extensions
of $O_Y$ contains precisely the permutations $\pi \in S_Y$ such that if $i \leq_{O_Y} j$, then $i$ precedes $j$ in $\pi$. Through the ordering of $v[Y]$, every permutation $\pi \in S_Y$ induces a canonical linear order on $v[Y]$, see [13]. Moreover, each permutation $\pi \in S_Y$ induces an acyclic orientation $O^\pi_Y \in \text{Acyc}(Y)$ defined by $O^\pi_Y(\{i, j\}) = (i, j)$ if $i$ precedes $j$ in $\pi$ and $O^\pi_Y(\{i, j\}) = (j, i)$ otherwise. The bijection

$$f_Y : S_Y / \sim_Y \rightarrow \text{Acyc}(Y), \quad f_Y([\pi]_Y) = O^\pi_Y,$$

from [13] allows us to identify equivalence classes and acyclic orientations. The number of equivalence classes under $\sim_Y$ is therefore given by $\alpha(Y)$.

For $O_Y \in \text{Acyc}(Y)$ and $e = \{v, w\} \in e[Y]$ let $O^e_Y$ be the orientation of $Y$ obtained from $O_Y$ by reversing the edge-orientation of $e$. Let $Y'_e$ and $Y''_e$ denote the graphs obtained from $Y$ by deletion and contraction of $e$, respectively, and let $O_{Y'_e}$ and $O_{Y''_e}$ denote the induced orientations of $O_Y$ under these operations. The bijection

$$\beta_e : \text{Acyc}(Y) \rightarrow \text{Acyc}(Y'_e) \cup \text{Acyc}(Y''_e)$$

defined by

$$O_Y \mapsto \begin{cases} O'_Y, & O^e_Y \notin \text{Acyc}(Y), \\ O'_Y, & O^e_Y \in \text{Acyc}(Y) \text{ and } O_Y(e) = (v, w), \\ O''_Y, & O^e_Y \in \text{Acyc}(Y) \text{ and } O_Y(e) = (w, v). \end{cases}$$

is well-known, and shows that one may compute $\alpha(Y)$ through the recursion relation

$$\alpha(Y) = \alpha(Y'_e) + \alpha(Y''_e),$$

valid for any $e \in e[Y]$.

## 3 Graph Constructions for Equivalence Relations

### 3.1 Relations on $S_Y / \sim_Y$

Using cycle notation, let $\sigma, \rho \in S_n$ be the elements

$$\sigma = (n, n-1, \ldots, 2, 1), \quad \rho = (1, n)(2, n-1) \cdots ([\frac{n}{2}], [\frac{n}{2}] + 1),$$

and let $C_n$ and $D_n$ be the subgroups

$$C_n = \langle \sigma \rangle \quad \text{and} \quad D_n = \langle \sigma, \rho \rangle.$$

Both $C_n$ and $D_n$ act on $S_Y$ via $g \cdot (\pi_1, \ldots, \pi_n) = (\pi_{g^{-1}(1)}, \ldots, \pi_{g^{-1}(n)})$. Define $\sigma_s(\pi) = \sigma^s \cdot \pi$, so that, e.g. $\sigma_1(\pi) = \sigma \cdot \pi = (\pi_2, \pi_3, \ldots, \pi_n, 1)$, and define $\rho(\pi) = \rho \cdot \pi = (\pi_n, \pi_{n-1}, \ldots, \pi_2, \pi_1)$. We construct two undirected graphs $C(Y)$ and $D(Y)$ whose vertex sets are $S_Y / \sim_Y$, and edge sets are

$$e[C(Y)] = \{[[\pi]_Y, [\sigma_1(\pi)]_Y] \mid \pi \in S_Y\},$$

$$e[D(Y)] = \{[[\pi]_Y, [\rho(\pi)]_Y] \mid \pi \in S_Y\} \cup e[C(Y)].$$
Define \(\kappa(Y)\) and \(\delta(Y)\) to be the number of connected components of \(C(Y)\) and \(D(Y)\), respectively. By construction, \(C(Y)\) is a subgraph of \(D(Y)\) and \(\delta(Y) \leq \kappa(Y)\).

**Example 1** Let \(Y\) be the complete bipartite graph \(K_{2,3}\), where the partition of the vertex set is \(\{\{1, 3, 5\}, \{2, 4\}\}\). The graph \(U(K_{2,3})\) is shown in Figure 1 with vertex labels omitted. By simply counting the components we see that \(\alpha(K_{2,3}) = 46\). We can better understand the component structure of \(U(K_{2,3})\) by mapping permutations as \((\pi_i) \mapsto (\pi_i \mod 2)\). Non-adjacency in \(Y\) coincides with parity, that is, if \(\pi \sim_Y \sigma\), then \(\phi(\pi) = \phi(\sigma)\). Through the map \(\phi\) we see that the 12 singleton points in \(U(K_{2,3})\) are precisely those with image 10101. Each of the 24 size-two components correspond to a pair of permutations with \(\phi\)-image of the form 01011, 11010, 01101, or 10110. The six square-components arise from the permutations with \(\phi\)-image 10011 and 11001. Finally, the permutations in the two hexagon-components are of the form 11100 or 00111.

The graphs \(C(K_{2,3})\) and \(D(K_{2,3})\) are shown in Figure 2. The dashed lines are edges in \(D(K_{2,3})\) but not in \(C(K_{2,3})\). The vertices in Figure 2 are labeled by a permutation in the corresponding equivalence class.
class in $S_Y / \sim_Y$. There are three isomorphic copies of the component on the right, but only one is shown. Each of these three components contains permutations whose $\phi$-image is in \{01101, 11010, 10101, 01011, 10110\}. The component on the left contains all of the remaining permutations, i.e., all $\pi$ for which $\phi(\pi) \in \{11100, 11001, 10011, 00111, 01110\}$. Clearly, $\kappa(K_{2,3}) = 7$ and $\delta(K_{2,3}) = 4$.

Permutations from $\sim_Y$ classes belonging to the same component in $C(Y)$ are called $\kappa$-equivalent permutations, as are the corresponding acyclic orientations. For two $\kappa$-equivalent permutations $\pi$ and $\pi'$ there is a sequence of adjacent non-edge transpositions and cyclic shifts that map $\pi$ to $\pi'$. This is simply a consequence of the definitions of $S_Y / \sim_Y$ and $C(Y)$. Similarly, two permutations belonging to $\sim_Y$ classes on the same connected component in $D(Y)$ are called $\delta$-equivalent, as are their corresponding acyclic orientations.

3.2 Relations on $\text{Acyc}(Y)$

The bijection between $S_Y / \sim_Y$ and $\text{Acyc}(Y)$ in (2) allows us to identify $[\pi]_Y$ with the acyclic orientation $O^\pi_Y$. It is clear that mapping $\pi \in [\pi']_Y$ to $\sigma_1(\pi)$ corresponds precisely to converting the vertex $\pi_1$ from a source to a sink in $O^\pi_Y$. This can be extended. Following [12] we call the conversion of a source vertex to a sink vertex in $O_Y \in \text{Acyc}(Y)$ a source-to-sink operation, or a click. Two orientations $O_Y, O'_Y \in \text{Acyc}(Y)$ where $O_Y$ can be transformed into $O'_Y$ by a sequence of clicks are said to be click-related, and we write this as $c(O_Y) = O'_Y$ where $c = c_1c_2 \cdots c_k$ with $c_i \in v[Y]$. It is straightforward to verify that this click-relation is an equivalence relation on $\text{Acyc}(Y)$. Through the bijection in (2) it is clear that a source-to-sink operation precisely encodes adjacency in the graph $C(Y)$, and the number of click-equivalence classes in $\text{Acyc}(Y)$ therefore equals $\kappa(Y)$. The second equivalence relation on $\text{Acyc}(Y)$ arises in the same manner by additionally identifying $O^\pi_Y$ and the reverse orientation $O^\pi_Y$, the unique orientation that satisfies $O^\pi_Y([i,j]) \neq O^\pi_Y([i,j])$ for every $\{i,j\} \in e[Y]$.

4 Structure of $C(Y)$ and $D(Y)$

The following result gives insight into the component structure of the graph $C(Y)$.

**Proposition 1** Let $Y$ be a connected graph on $n$ vertices and let $g, g' \in C_n$ with $g \neq g'$. Then $[g \cdot \pi]_Y \neq [g' \cdot \pi]_Y$.

**Proof** Assume $g \neq g'$ with $[g \cdot \pi]_Y = [g' \cdot \pi]_Y$. By construction, we have $g \cdot \pi = \sigma_s(\pi)$ and $g' \cdot \pi = \sigma_{s'}(\pi)$. Without loss of generality we may assume $s' < s$. Let $V' \subset V = v[Y]$ be the initial subsequence of vertices in $\sigma_{s'}(\pi)$ that occurs at the end in $\sigma_s(\pi)$. If any of the vertices in $V'$ are adjacent to any of the vertices in $V \setminus V'$ in $Y$ it would imply that
\[ [\sigma_s(\pi)]_Y \neq [\sigma_{s'}(\pi)]_Y. \] The only possibility is that \( Y \) is not connected, but this contradicts the assumptions of the proposition. \( \square \)

There is a similar result to Proposition 1 for \( D_n \), albeit somewhat more restrictive.

**Proposition 2** Let \( Y \) be a connected graph on \( n \) vertices and let \( g, g' \in D_n \) with \( g \neq g' \). If \( [g \cdot \pi]_Y = [g' \cdot \pi]_Y \) holds then \( Y \) must be bipartite.

**Proof** From \( [g \cdot \pi]_Y = [g' \cdot \pi]_Y \) it follows from Proposition 1 that \( g \) and \( g' \) lie in different cosets of \( C_n \) in \( D_n \). Without loss of generality we may assume that \( g = \sigma^s \) and \( g' = \rho \sigma^{s'} \). Let \( m = |s' - s| \) and \( n' = n - m \). If \( s' > s \) (resp. \( s' < s \)) the first (resp. last) \( m \) elements of \( g \cdot \pi \) and \( g' \cdot \pi \) are the same but occur in reverse order. Call the set of these elements \( V_1 \). The remaining \( n' \) elements occur in reverse order as well in the two permutations. Let \( V_2 \) denote the set of these elements. For \( [g \cdot \pi]_Y = [g' \cdot \pi]_Y \) to hold, there cannot be an edge between any two vertices in \( V_1 \), or between any two vertices in \( V_2 \). Therefore, the graph \( Y \) must be a subgraph of \( K(V_1, V_2) \), the complete bipartite graph with vertex sets \( V_1 \) and \( V_2 \). \( \square \)

**Remark 1** The pairs \((\sigma^s, \rho \sigma^{s'})\) and \((\sigma^{s'}, \rho \sigma^s)\) determine the same bipartite graph in the proof above. Also, the vertex sets \( V_1 \) and \( V_2 \) can only consist of consecutive elements in \( \pi \).

**Remark 2** If \( Y \) is connected and bipartite then \(|\{[g \cdot \pi]_Y \mid g \in D_n\}| = 2n - 1 \). This follows from the fact that at most two \( \sim_{\pi} \) classes can coincide as all distinct pairs \( g \) and \( g' \) for which equality holds leads to different sets \( V_1(\{g, g'\}) \) and \( V_2(\{g, g'\}) \) modulo Remark 1. The existence of two or more distinct partitions of \( v[Y] \) into sets \( V_1 \) and \( V_2 \) as above would imply that \( Y \) is not connected.

We next consider the quantity \( \delta(Y) \), and will show how it is determined by \( \kappa(Y) \).

**Lemma 1** The reflection map \( \rho : S_Y \longrightarrow S_Y \) extends to an involution

\[
\rho^* : \text{Acyc}(Y)/\sim_{\kappa} \longrightarrow \text{Acyc}(Y)/\sim_{\kappa}.
\]

**Proof** By the definition of \( U(Y) \) it follows that if \( \pi, \pi' \in S_Y \) are adjacent in \( U(Y) \) then so are \( \rho(\pi) \) and \( \rho(\pi') \). It follows easily that \( \pi \sim_{\pi'} \) implies \( \rho(\pi) \sim_{\pi'} \rho(\pi') \). The map \( \rho \) therefore extends to a map \( \hat{\rho} : S_Y/\sim_{\pi} \longrightarrow S_Y/\sim_{\pi'} \) by \( \hat{\rho}(\pi'|_Y) = [\rho(\pi)|_Y \]. Likewise, if \( O_Y \) and \( O_{Y'} \) are \( \kappa \)-equivalent then so are \( \hat{\rho}(O_Y) \) and \( \hat{\rho}(O_Y) \) [using the bijection (2)], and \( \hat{\rho} \) extends to \( \rho^* \) as in (3) by \( \rho^*(A) = \hat{\rho}(O_A) \) for any \( O_A \in \text{Acyc}(Y)/\sim_{\kappa} \). This map is clearly an involution since \( \rho \) itself is an involution. \( \square \)

**Proposition 3** Let \( Y \) be a connected undirected graph. If \( Y \) is not bipartite then \( \delta(Y) = \frac{1}{2} \kappa(Y) \). If \( Y \) is bipartite then \( \delta(Y) = \frac{1}{2}(\kappa(Y) + 1) \).
Proof If $Y$ is not bipartite then Proposition \ref{prop:involution} ensures that the involution $\rho^*$ has no fixed points, from which we conclude that $2\delta(Y) = \kappa(Y)$. For the second statement, we use Proposition \ref{prop:equivalence} and Remark \ref{rem:fixed_point} and the connectedness of $Y$ to conclude that $\rho^*$ has precisely one fixed point. Since $\rho$ is an involution it follows that $\delta(Y) = \kappa(Y) - 1 = \frac{\kappa(Y) + 1}{2} = \left\lceil \frac{\kappa(Y)}{2} \right\rceil$.

Thus, we always have $\delta(Y) = \left\lceil \frac{\kappa(Y)}{2} \right\rceil$, and we also have the following characterization of bipartite graphs:

**Corollary 1** A connected graph $Y$ is bipartite if and only if $\kappa(Y)$ is odd.

For an example where $\rho^*$ has a fixed point see Figure 2 in Example \ref{ex:fixed_point}. From Proposition \ref{prop:equivalence} we can derive an upper bound for $\kappa(Y)$.

**Proposition 4** If $Y$ is a connected undirected graph on $n$ vertices, then $\kappa(Y) \leq \frac{1}{n} \alpha(Y)$.

Proof By Proposition \ref{prop:equivalence} for any $\pi \in S_Y$, the set \{O_{\sigma_1}(\pi)Y, \ldots, O_{\sigma_n}(\pi)Y\} contains $n$ distinct acyclic orientations that are all $\kappa$-equivalent and the proof follows. □

This bound is sharp for certain graphs such as the complete graph $K_n$.

5 Poset Structure of $\kappa$-Equivalence Classes

An edge $e$ of an undirected graph $Y$ is a bridge if removing $e$ increases the number of connected components of $Y$. An edge that is not a bridge is a cycle-edge, or equivalently, an edge $e$ is a cycle-edge if it is contained in a cycle traversing $e$ precisely once. The main result in \cite{9} is a recurrence relation for $\kappa(Y)$ under edge deletion and edge contraction.

**Theorem 1** (\cite{9}) Let $Y$ be a finite undirected graph with $e \in e[Y]$, and let $Y'_e$ be the graph obtained from $Y$ by deleting $e$, and let $Y''_e$ be the graph obtained from $Y$ by contracting $e$. Then

$$\kappa(Y) = \begin{cases} \kappa(Y_1)\kappa(Y_2), & e \text{ is a bridge linking } Y_1 \text{ and } Y_2, \\ \kappa(Y'_e) + \kappa(Y''_e), & e \text{ is a cycle-edge}. \end{cases}$$ (6)

The first case implies that $\kappa(Y)$ is unaffected by removal of bridges, and is relatively straightforward to establish. However, the case where $e$ is a cycle-edge is harder, and was proven in \cite{3} on the level of acyclic orientations. In this section, we will show how one may associate a poset with each $\kappa$-equivalence class. The properties of this poset give us better insight into the structure of $\text{Acyc}(Y)/\sim_\kappa$. Additionally, it allows us to construct an alternative proof for Theorem \ref{thm:equivalence}.

Throughout, we will let $e = \{v, w\}$ be a fixed cycle-edge of the connected graph $Y$, and for ease of notation we set $Y' = Y'_e$ and $Y'' = Y''_e$. For $O \in \text{Acyc}(Y)$ we let $O_Y$, and $O_Y''$ denote the induced orientations of $Y'$. 


and \(Y''\). Notice that \(O_Y\) is always acyclic, while \(O_{Y''}\) is acyclic if and only if there is no directed path with endpoints \(v\) and \(w\) in \(O_Y\). Finally, we let \([O_Y]\) denote the \(\kappa\)-equivalence class containing \(O_Y\).

The interval \([a, b]\) of a poset \(P\) (where \(a \leq b\)) is the subposet consisting of all \(c \in P\) such that \(a \leq c \leq b\). Viewing a finite poset \(P\) as a directed graph \(D_P\), the interval \([a, b]\) contains precisely the vertices that lie on a directed path from \(a\) to \(b\), and thus is a vertex-induced subgraph of \(D_P\).

By assumption, \(Y\) contains the edge \(\{v, w\}\), so for all \(O_Y \in \text{Acyc}(Y)\) either \(v \leq_{O_Y} w\) or \(w \leq_{O_Y} v\). In this section, we will study the interval \([v, w]\) in the poset \(O_Y\) (when \(v \leq_{O_Y} w\)) and its behavior under clicks.

**Definition 1** Let \(\text{Acyc}_<(Y)\) be the set of acyclic orientations of vertex-induced subgraphs of \(Y\). We define the map

\[
I : \text{Acyc}(Y) \longrightarrow \text{Acyc}_<(Y),
\]

by \(I(O_Y) = [v, w]\) if \(v \leq_{O_Y} w\), and by \(I(O_Y) = \emptyset\) (the null graph) otherwise. When \(I(O_Y) \neq \emptyset\) we refer to \(I(O_Y)\) as the \(vw\)-interval of \(O_Y\).

Elements of \(\text{Acyc}_<(Y)\) can be thought of as subposets of \(\text{Acyc}(Y)\). Through a slight abuse of notation, we will at times refer to \(I(O_Y)\) as a poset, a directed graph, or a subset of \(v[O_Y]\). In this last case, it is understood that the relations are inherited from \(O_Y\).

For an undirected path \(P = v_1, v_2, \ldots, v_k\) in \(Y\), we define the function

\[
\nu_P : \text{Acyc}(Y) \longrightarrow \mathbb{Z},
\]

where \(\nu_P(O_Y)\) is the number of edges oriented as \((v_i, v_{i+1})\) in \(O_Y\), minus the number of edges oriented as \((v_{i+1}, v_i)\). It is clear that if \(P\) is a cycle, then \(\nu_P\) is preserved under clicks, and in this case, \(\nu_P\) extends to a map

\[
\nu_P : \text{Acyc}(Y) / \sim_\kappa \longrightarrow \mathbb{Z}.
\]

We will now prove a series of structural results about the \(vw\)-interval. Since \(\{v, w\} \in e[Y]\), every \(\kappa\)-class contains at least one orientation \(O_Y\) with \(v \leq_{O_Y} w\), and thus there is at least one element \(O_Y\) in each \(\kappa\)-class with \(I(O_Y) \neq \emptyset\). The next results shows how we can extend the notion of \(vw\)-interval from over \(\text{Acyc}(Y)\) to \(\text{Acyc}(Y) / \sim_\kappa\).

**Proposition 5** The map \(I\) can be extended to a map

\[
I^* : \text{Acyc}(Y) / \sim_\kappa \longrightarrow \text{Acyc}_<(Y)
\]

by \(I^*([O_Y]) = I(O_Y)\), where \(O_Y^1\) is any element of \([O_Y]\) for which \(I(O_Y^1) \neq \emptyset\).

**Proof** It suffices to prove that \(I^*\) is well-defined. Consider \(O_Y^1 \sim_\kappa O_Y^2\), with \(v \leq_{O_Y}^i w\) for \(i = 1, 2\). To show that \(I(O_Y^1) = I(O_Y^2)\) let \(a\) be a vertex in \(I(O_Y^1)\). Then \(a\) lies on a directed path \(P'\) from \(v\) to \(w\) in \(O_Y^1\), say of length \(k \geq 2\). Let \(P\) be the cycle formed by adding vertex \(v\) to the end of \(P'\). Clearly \(\nu_P(O_Y^1) = k - 1\) since \(O_Y^1(v) = (v, w)\).
By assumption, $O_Y^2 \in [O_Y^1]$ with $v \leq_{O_Y^2} w$. Since $\nu_P$ is constant on $[O_Y^1]$ it follows from $\nu_P(O_Y^1) = k - 1 = \nu_P(O_Y^2)$ that every edge of $P'$ is oriented identically in $O_Y^1$ and $O_Y^2$, and hence that every directed path $P'$ in $O_Y^1$ is contained in $O_Y^2$ as well. The reverse inclusion follows by an identical argument. □

In light of Proposition 5 we define the $vw$-interval of a $\kappa$-class $[O_Y]$ to be $\mathcal{I}^*(\{O_Y\})$. The $vw$-interval will be central in understanding properties of click-sequences. First, we will make a simple observation without proof, which also appears in [10] in the context of admissible sequences in Coxeter theory.

**Proposition 6** Let $O_Y \in \text{Acyc}(Y)$, let $c = c_1c_2 \cdots c_m$ be an associated click-sequence, and consider any directed edge $(v_1,v_2)$ in $O_Y$. Then the occurrences of $v_1$ and $v_2$ in $c$ alternate, with $v_1$ appearing first.

Because $\{v,w\} \in e[Y]$, we can say more about the vertices in $\mathcal{I}(O_Y)$ that appear between successive instances in $v$ and $w$ in a click-sequence.

**Proposition 7** Let $O_Y \in \text{Acyc}(Y)$, and let $c = c_1c_2 \cdots c_m$ be an associated click-sequence that contain every vertex of $\mathcal{I}(O_Y)$ at least once and with $c_1 = v$. Then every vertex of $\mathcal{I}(O_Y)$ appears in $c$ before any vertex in $\mathcal{I}(O_Y)$ appears twice.

**Proof** The proof is by contradiction. Assume the statement is false, and let $a \in \mathcal{I}(O_Y)$ be the first vertex whose second instance in $c$ occurs before the first instance of some other vertex $z \in \mathcal{I}(O_Y)$. If $a \neq v$, then $a$ is not a source in $O_Y$, and there exists a directed edge $(a',a)$. By Proposition 6 $a'$ must appear in $c$ before the first instance of $a$, but also between the two first instances of $a$. This is impossible, because $a$ was chosen to be the first vertex appearing twice in $c$. That only leaves $a = v$, and $v$ must appear twice before the first instance of $w$. However, this contradicts the statement of Proposition 6 because $\{v,w\} \in e[Y]$. □

The next result shows that for any click-sequence $c$ that contains every element in $\mathcal{I}(O_Y)$ precisely once, we may assume without loss of generality that the vertices in $\mathcal{I}(O_Y)$ appear consecutively.

**Proposition 8** Let $O_Y \in \text{Acyc}(Y)$ be an acyclic orientation with $v \leq_{O_Y} w$. If $c = c_1c_2 \cdots c_m$ is an associated click-sequence containing precisely one instance of $w$, and no subsequent instances of vertices from $\mathcal{I}(O_Y)$, then there exists a click-sequence $c' = c'_1c'_2 \cdots c'_m$ such that (i) there exists an interval $[p,q]$ of $\mathbb{N}$ with $c'_j \in \mathcal{I}(O_Y)$ iff $p \leq j \leq q$, and (ii) $c(O_Y) = c'(O_Y)$.

**Proof** We prove the proposition by constructing a desired click-sequence $c''$ from $c$ through a series of transpositions where each intermediate click-sequence $c'$ satisfies $c(O_Y) = c'(O_Y)$. Such transpositions are said to have property $T$.

Let $I = \mathcal{I}(O_Y)$, and let $A$ be the set of vertices in $I^c = v[Y] \setminus I$ that lie on a directed path in $O_Y$ to a vertex in $I$ (vertices above $I$), and let $B$ be the
set of vertices that lie on a directed path in \( O_Y \) from a vertex in \( I \) (vertices below \( I \)). Let \( C \) be the complement of \( I \cup A \cup B \). Two vertices \( c_i, c_j \in A \cup B \) with \( i < j \) for which there is no element \( c_k \in A \cup B \) with \( i < k < j \) are said to be \textit{tight}. We will investigate when transpositions of tight vertices in a click-sequence \( \mathbf{c} \) of \( O_Y \) has property \( T \), and we will see that this is always the case if \( c_i \in B \) and \( c_j \in A \). Consider the intermediate acyclic orientation after applying successive clicks \( c_1 c_2 \cdots c_{i-1} \) to \( O_Y \). Obviously, \( c_i \) is a source. At this point, if \( c_j \) were not a source, then there would be an adjacent vertex \( a \in A \) with the edge \( \{a, c_j\} \) oriented \( (a, c_j) \). For \( c_j \) to be clicked as usual (i.e., as a source), \( a \) must be clicked first, but this would break the assumption that \( c_i \) and \( c_j \) are tight. Therefore, \( c_i \) and \( c_j \) are both sources at this intermediate step, and so the vertices \( c_i, c_{i+1}, \ldots, c_j \) are an independent set of sources, and may be permuted in any manner without changing the image of the click sequence. Therefore, the transposition of \( c_i \) and \( c_j \) in \( \mathbf{c} \) has property \( T \), as claimed. By iteratively transposing tight pairs in \( \mathbf{c} \), we can construct a click-sequence with the property that every vertex in \( A \) comes before every vertex in \( B \). In light of this, we may assume without loss of generality that \( \mathbf{c} \) has this property.

The next step is to show that we can move all vertices in \( A \) before \( v \), and all vertices in \( B \) after \( w \) via transpositions having property \( T \). Let \( a \) be the first vertex in \( A \) appearing after \( v \) in the click sequence \( \mathbf{c} \). We claim that the transposition moving \( a \) to the position directly preceding \( v \) has property \( T \). This is immediate from the observation that when \( v \) is to be clicked, \( a \) is a source as well, by the definition of \( A \), thus it may be clicked before \( v \), without preventing subsequent clicks of vertices up until the original position of \( a \). Therefore, we may one-by-one move the vertices in \( A \) that are between \( v \) and \( w \), in front of \( v \). An analogous argument shows that we may move the vertices in \( B \) that appear before \( w \) to a position directly following \( w \). In the resulting click-sequence \( \mathbf{c}' \), the only vertices between \( v \) and \( w \) are either in \( I \) or \( C \). The subgraph of the directed graph \( O_Y \) induced by \( C \) is a disjoint union of weakly connected components, and none of the vertices are adjacent to \( I \). By definition of \( A \) and \( B \), there cannot exist a directed edge \( (c, a) \) or \( (b, c) \), where \( a \in A \), \( b \in B \), and \( c \in C \). Thus for each weakly connected component of \( C \), the vertices in the component can be moved within \( \mathbf{c}' \), preserving their relative order, to a position either (i) directly after the vertices in \( A \) before \( v \), or (ii) directly after \( w \) and before the vertices of \( B \). Call this resulting click-sequence \( \mathbf{c}'' \). As we just argued, all the transpositions occurring in the rearrangement \( \mathbf{c} \mapsto \mathbf{c}'' \) has property \( T \), and \( \mathbf{c}'' \) contains all of the vertices in \( I \) in consecutive order, and this proves the result. \( \Box \)

We remark that the last two results together imply that for the interval \([p, q]\) in the statement of Proposition \( \square \), \( c_p = v \), \( c_q = w \), and the sequence \( c_p c_{p+1} \cdots c_q \) contains every vertex in \( I(O_Y) \) precisely once. A simple induction argument implies the following.
The argument is by induction on a click-sequence $c$. Suppose the statement holds for all intervals $[p_i, q_i]$ of $\mathbb{N}$ such that $c_j \in \mathcal{I}(O_Y)$ iff $p_i \leq j \leq q_i$ for some $i$, and (ii) $c(O_Y) = c'(O_Y)$.

Proof The argument is by induction on $k$. When $k = 1$, the statement is simply Proposition 8. Suppose the statement holds for all $k \leq N$, for some $N \in \mathbb{N}$, and let $c$ be a click-sequence containing $N + 1$ instances of $w$. Let $c^\ell$ be the second instance of $v$ in $c$, and consider the two click-sequences $c_i := c_1c_2 \cdots c_{\ell-1}$ and $c_f := c_{\ell}c_{\ell+1} \cdots c_m$. By Proposition 8 there exists an interval $[p_1, q_1]$ with $p_1 < q_1 < \ell$, and by the induction hypothesis, there exists $k$ intervals $[p_2, q_2], \ldots, [p_k, q_k]$ with $\ell \leq p_2 < q_2 < \cdots < p_k < q_k < q_{k+1}$ such that if $c_j \in \mathcal{I}(O_Y)$, then $p_i \leq j \leq q_i$ for some $i = 1, \ldots, k + 1$.

Let $\eta_k : \text{Acyc}(Y) \hookrightarrow \text{Acyc}(Y')$ be the restriction map that sends $O_Y$ to $O_{Y'}$. Clearly, this map extends to a map $\eta_k^* : \text{Acyc}(Y)/\sim_\kappa \rightarrow \text{Acyc}(Y')/\sim_\kappa$. Define

$$\mathcal{I}_\kappa^*: \text{Acyc}(Y')/\sim_\kappa \rightarrow \text{Acyc}_{\leq}(Y')$$

by $\mathcal{I}_\kappa^*(\{O_Y\}) = \mathcal{I}(O_Y)$ for any $O_Y \in [O_Y]$ such that $\eta_k([O_Y]) = [O_Y]$ with $|\mathcal{I}(O_Y)| \geq 3$, and $\mathcal{I}_\kappa^*(\{O_Y\}) = \{v, w\}$ if no such acyclic orientation $O_Y$ exists.

**Proposition 9** The map $\mathcal{I}_\kappa^*$ is well-defined, and the diagram

$$
\begin{array}{ccc}
\text{Acyc}(Y)/\sim_\kappa & \xrightarrow{\eta_k^*} & \text{Acyc}_{\leq}(Y) \\
\downarrow & & \downarrow \\
\text{Acyc}(Y')/\sim_\kappa & \xrightarrow{\mathcal{I}_\kappa^*} & \text{Acyc}_{\leq}(Y')
\end{array}
$$

commutes.

Proof Let $[O_Y] \in \text{Acyc}(Y')/\sim_\kappa$. If there is at most one orientation $O_Y \in \text{Acyc}(Y)$ such that $|\mathcal{I}(O_Y)| \geq 3$ and $\eta_k(O_Y) \in [O_Y']$, or if all orientations of the form $O_Y$ in the definition of $\mathcal{I}_\kappa^*$ are $\kappa$-equivalent, then both statements of the proposition are clear. Assume therefore that there are acyclic orientations $O^*_Y, O'_Y \in \text{Acyc}(Y)$ with $O^*_Y \sim_\kappa O'_Y$, but $\eta_k^*([O^*_Y]) = \eta_k^*([O'_Y])$ and $|\mathcal{I}(O^*_Y)|, |\mathcal{I}(O'_Y)| \geq 3$. It suffices to prove that in this case,

$$\mathcal{I}(O^*_Y) = \mathcal{I}(O'_Y). \quad (7)$$

This is equivalent to showing that the set of $vw$-paths (directed paths from $v$ to $w$) in $O^*_Y$, is the same as the set of $vw$-paths in $O'_Y$. From this it will also follow that the diagram commutes. By assumption, both of these orientations contain at least one $vw$-path. We will consider separately the cases when these orientations share or do not share a common $vw$-path.

Case 1: $O^*_Y$ and $O'_Y$ share no common $vw$-path. Let $P_1$ be a length-$k_1$ $vw$-path in $O^*_Y$, and let $P_2$ be a length-$k_2$ $vw$-path in $O'_Y$. Suppose
Moreover, if \( \pi \) to a single vertex denoted \( \mathcal{I}(O_Y) \) has at least two vertices. We write \( Y_I \) for the graph formed from \( Y \) by contracting all vertices in \( I \) to a single vertex denoted \( V_I \). If \( I \) only contains \( v \) and \( w \) then \( Y_I = Y_v'' \). Moreover, \( O_Y \) gives rise to an orientation \( O_{Y_I} \) of \( Y_I \), and this orientation is clearly acyclic.

**Proposition 10** Let \( O_Y^1, O_Y^2 \in \text{Acyc}(Y) \) and assume \( \mathcal{I}(O_Y^1) = \mathcal{I}(O_Y^2) \). If \( O_Y^1 \sim \kappa, O_Y^2 \sim \kappa \) then \( [O_Y^1] \sim_{\kappa} [O_Y^2] \).

**Proof** We prove the contrapositive statement. Set \( I = \mathcal{I}(O_Y^1) \), suppose \( |I| = k \), and let \( v_1 v_2 \cdots v_k \) be a linear extension of \( I \). For any click-sequence \( c_I \) between two acyclic orientations \( O_Y^1 \) and \( O_Y^2 \) in \( \text{Acyc}(Y_I) \), let \( c \) be the click-sequence formed by replacing every occurrence of \( c_i = V_i \) in \( c_I \) by the sequence \( v_1 \cdots v_k \). Then \( c(O_Y^1) = O_Y^2 \) and \( O_Y^1 \sim \kappa, O_Y^2 \sim \kappa \) as claimed. \( \square \)

### 6 Proof of Theorem 1

In this section, we will utilize the results in the previous section to establish a bijection from \( \text{Acyc}(Y/\sim_\kappa) \) to \( (\text{Acyc}(Y''/\sim_\kappa) \cup \text{Acyc}(Y''''/\sim_\kappa) \) for any cycle-edge \( e \), which will in turn imply Theorem 1.

For \([O_Y] \in \text{Acyc}(Y/\sim_\kappa)\), let \( O_Y^e \) denote an orientation in \([O_Y]\) such that \( \pi = v \pi_2 \cdots \pi_n \) and \( w = \pi_i \) for \( i \) minimal. Define the map

\[
\Theta: \text{Acyc}(Y/\sim_\kappa) \rightarrow (\text{Acyc}(Y'/\sim_\kappa) \cup (\text{Acyc}(Y''/\sim_\kappa) \cup (\text{Acyc}(Y''''/\sim_\kappa)) \]

by

\[
[O_Y] \mapsto \begin{cases} [O_Y^e], & \exists O_Y^e \in [O_Y] \text{ with } \pi = vw \pi_3 \cdots \pi_n \\ [O_Y^e], & \text{otherwise.} \end{cases}
\]

where \( O_Y^e \) is the orientation in \([O_Y]\) such that \( \pi = v \pi_2 \cdots \pi_n \) and \( w = \pi_i \) for \( i \) minimal. Define the map

\[
\Theta: \text{Acyc}(Y/\sim_\kappa) \rightarrow (\text{Acyc}(Y'/\sim_\kappa) \cup (\text{Acyc}(Y''/\sim_\kappa) \cup (\text{Acyc}(Y''''/\sim_\kappa)) \]

by

\[
[O_Y] \mapsto \begin{cases} [O_Y^e], & \exists O_Y^e \in [O_Y] \text{ with } \pi = vw \pi_3 \cdots \pi_n \\ [O_Y^e], & \text{otherwise.} \end{cases}
\]

and \( O_Y^e \) is the orientation in \([O_Y]\) such that \( \pi = v \pi_2 \cdots \pi_n \) and \( w = \pi_i \) for \( i \) minimal.
Note that $[O_Y]$ is mapped into $\text{Acyc}(Y''')/\sim_κ$ if and only if the only vertices in $I^∗(O_Y)$ are $v$ and $w$. Since $κ$-equivalence over $Y$ implies $κ$-equivalence over $Y'$, $Θ$ does not depend on the choice of $π$, and is thus well-defined. The results we have derived for the $vw$-interval now allow us to establish the following:

**Theorem 2** The map $Θ$ is a bijection.

**Proof** We first prove that $Θ$ is surjective. Let $I = \{v, w\}$ and consider an element $[O_Y'''] ∈ \text{Acyc}(Y''')/\sim_κ$ with $O_Y''' ∈ [O_Y''']$ where $π = V_1π_2⋯π_{n−1}$. Let $π^+ = vwπ_2⋯π_{n−1} ∈ S_Y′$. Clearly $[O_Y^{π^+}] ∈ \text{Acyc}(Y)/\sim_κ$ is mapped to $[O_Y]$, by $Θ$.

Next, consider an element $[O_Y] ∈ \text{Acyc}(Y')/\sim_κ$. If there is no element $O_Y^π$ of $[O_Y']$ such that $π = vwπ_3⋯π_{n}$, then no elements of $[O_Y]$ are of this form either, and by definition $[O_Y]$ has a preimage under $Θ$. We are left with the case where $[O_Y']$ contains an element $O_Y^π$, such that $π = vwπ_3⋯π_{n}$, and we must show that there exists $O_Y^{π'} ∈ [O_Y']$ such that $[O_Y^{π'}]$ contains no element of the form $O_Y^π$ with $σ = vwσ_3⋯σ_{n}$. Note that if $σ = vwσ_3⋯σ_{n}$, then the vertices in $I(O_Y^{π'})$ are precisely $v$ and $w$. If the orientation $O_Y^{π'}$ had a directed path from $v$ to $w$, then the corresponding orientation $O_Y ∈ \text{Acyc}(Y)$ formed by adding the edge $e$ with orientation $(v, w)$ has $vw$-interval of size at least $3$, so by Proposition 5 the acyclic orientation $O_Y$ cannot be $κ$-equivalent to any orientation $O_Y^σ$ such that $σ = vwσ_3⋯σ_{n}$.

Thus it remains to consider the case when $[O_Y]$ contains no acyclic orientation with a directed path from $v$ to $w$. Pick any simple undirected path $P'$ from $v$ to $w$ in $Y'$, which exists because $e$ is a cycle-edge. Choose an orientation in $[O_Y']$ for which $ν_{P'}$ is maximal. Without loss of generality we may assume that $O_Y′$ is this orientation. Let $O_Y ∈ \text{Acyc}(Y)$ be the orientation that agrees with $O_Y′$, and with $e$ oriented as $(w, v)$. Since we have assumed that there is no directed path from $v$ to $w$ this orientation is acyclic. We claim that for any $σ = vwσ_3⋯σ_{n}$ one has $O_Y^σ / [O_Y']$. To see this, assume the statement is false. Let $P$ be the undirected cycle in $Y$ formed by adding the edge $e$ to the path $P′$. Because $e$ is oriented as $(v, w)$ in $O_Y^σ$ and as $(w, v)$ in $O_Y$, we have $ν_{P'}(O_Y^σ) − ν_{P'}(O_Y) = ν_{P'}(O_Y) + 1$. If $O_Y$ and $O_Y^σ$ were $κ$-equivalent, then

$$ν_{P'}(O_Y^σ) − 1 = ν_{P'}(O_Y^σ) = ν_{P'}(O_Y) − 1,$$

and thus $ν_{P'}(O_Y^σ) = ν_{P'}(O_Y) + 2$. Any click sequence mapping $O_Y$ to $O_Y^σ$ is a click-sequence from $O_Y′$, to $O_Y^σ$. Therefore, $O_Y^σ ∈ [O_Y']$, which contradicts the maximality of $ν_{P'}(O_Y′)$. We therefore conclude that $O_Y^σ / [O_Y']$, that $Θ([O_Y]) = [O_Y]$, and hence that $Θ$ is surjective.

We next prove that $Θ$ is an injection. By Proposition 11 (with $I = \{v, w\}$), $Θ$ is injective when restricted to the preimage of $[O_Y']$ under $Θ$. Thus it suffices to show that every element in $\text{Acyc}(Y')/\sim_κ$ has a unique preimage under $Θ$. By Proposition 8 every preimage of $[O_Y']$ must have the same $vw$-interval $I$, containing $k > 2$ vertices. Suppose there were preimages
when restricted to their sets of periodic points. By Proposition \[ \text{(10)} \] it follows that \( O^\pi_{Y_j} \sim_k O^\pi_{Y_1} \). We will now show that this leads to a contradiction.

Assume that \( c = c_1 \cdots c_m \) is a click-sequence from \( O^\pi_{Y_j} \) to \( O^\pi_{Y_r} \). If one of \( \pi \) or \( \sigma \) is not \( \kappa \)-equivalent to a permutation with vertices \( v \) and \( w \) in succession, then their corresponding \( \kappa \)-classes would be unchanged by the removal of edge \( e \). In light of this, we may assume that \( \pi = v\pi_2 \cdots \pi_{n-1}w \) and \( \sigma = v\sigma_2 \cdots \sigma_{n-1}w \), and thus that \( c_1 = v \) and \( c_m = w \). By Proposition \[ \text{(8)} \] we may assume that the vertices in \( I \) appear in \( c \) in some number of disjoint consecutive “blocks,” i.e., subsequences of the form \( c_i \cdots c_{i+k-1} \). Replacing each of these blocks with \( V_I \) yields a click-sequence from \( O^\pi_{Y_j} \) to \( O^\pi_{Y_1} \), contradicting the fact that \( O^\pi_{Y_j} \sim_k O^\pi_{Y_1} \). Therefore, no such click sequence \( c \) exists, and \( \Theta \) must be an injection, and the proof is complete. \( \square \)

Clearly, Theorem \[ \text{2} \] implies Theorem \[ \text{1} \] It is also interesting to note that the bijection \( \beta_c : \text{Acyc}(Y) \rightarrow \text{Acyc}(Y'_c) \cup \text{Acyc}(Y''_c) \) in \[ \text{(3)} \] does not extend to a map on \( \kappa \)-classes.

7 Discrete Dynamical Systems, Coxeter Groups, Node-firing Games, and Quiver Representations

We conclude with a brief discussion of how the equivalence relations studied in this paper arise in various areas of mathematics. The original motivation came from the authors’ interest in sequential dynamical systems (SDSs). The equivalence relation \( \sim_Y \) arises naturally in the study of functional equivalence of these systems. This can be seen as follows. Given a graph \( Y \) with vertex set \( \{1, 2, \ldots, n\} \) as above, a state \( x_v \in K \) is assigned to each vertex \( v \) of \( Y \) for some finite set \( K \). The system state is the tuple consisting of all the vertex states, and is denoted \( x = (x_1, \ldots, x_n) \in K^n \). The sequence of states associated to the 1-neighborhood \( B_1(v; Y) \) of \( v \) in \( Y \) (in some fixed order) is denoted \( x[v] \). A sequence of vertex functions \( (f_i)_j \) with \( f_i : K^{d(i)+1} \rightarrow K \) induces \( Y \)-local functions \( F_i : K^n \rightarrow K^n \) of the form

\[
F_i(x_1, \ldots, x_n) = (x_1, \ldots, x_{i-1}, f_i(x[i]), x_{i+1}, \ldots, x_n).
\]

The sequential dynamical system map with update order \( \pi = (\pi_i)_i \in S_Y \) is the function composition

\[
[\mathcal{S}_Y, \pi] = F_{\pi(n)} \circ F_{\pi(n-1)} \circ \cdots \circ F_{\pi(2)} \circ F_{\pi(1)}. \quad (10)
\]

By construction, if \( \pi \sim_Y \pi' \) holds then \( [\mathcal{S}_Y, \pi] \) and \( [\mathcal{S}_Y, \pi'] \) are identical as functions. Thus, \( \alpha(Y) \) is an upper bound for the number of functionally non-equivalent SDS maps that can be generated over the graph \( Y \) for a fixed sequence of vertex functions. For any graph \( Y \), there exist \( Y \)-local functions for which this bound is sharp \[ \text{[1]} \]. A weaker form of equivalence is cycle equivalence, which means that the dynamical system maps are conjugate (using the discrete topology) when restricted to their sets of periodic points.
In the language of graph theory, this means their periodic orbits are isomorphic as directed graphs. The following result shows how $\kappa$- and $\delta$-equivalent update orders yield dynamical system maps that are cycle equivalent.

**Theorem 3** For any finite set $K$ of vertex states, and for any $\pi \in S_Y$, the SDS maps $[F_Y, \pi]$ and $[F_Y, \sigma_s(\pi)]$ are cycle equivalent. For vertex states $K = F_2$ the SDS maps $[F_Y, \pi]$ and $[F_Y, \rho(\pi)]$ are cycle equivalent as well.

We refer to [8] for the proof of this result. Two sequential dynamical systems are cycle equivalent if their phase space digraphs are isomorphic when restricted to the cycles. The paper [8] contains additional background on equivalences of sequential dynamical systems as well as applications of $\kappa$-equivalence to the structural properties of their phase spaces.

There is also a connection between $\kappa$-equivalence and the theory of Coxeter groups. There is a natural bijection between the set of Coxeter elements $\prod_i s_{\pi(i)}$ of a Coxeter group (see, e.g. [1]) with generators $s_i$ for $1 \leq i \leq n$ and Coxeter graph $Y$ (ignoring order labels), and the set of equivalence classes $[\pi]_Y$. This is clear since the commuting generators are precisely those that are not connected in $Y$. Let $C(W, S)$ denote the set of Coxeter elements of a Coxeter group $W$ with generating set $S = \{s_1, \ldots, s_n\}$. It is shown in [14] that there is a bijection

$$C(W, S) \rightarrow \text{Acyc}(Y).$$  

Moreover, conjugating a Coxeter element $\prod_i s_{\pi(i)}$ by $s_{\pi(1)}$ corresponds to a cyclic shift, i.e.,

$$s_{\pi(1)}(s_{\pi(1)}s_{\pi(2)} \cdots s_{\pi(n)})s_{\pi(1)} = s_{\pi(2)} \cdots s_{\pi(n)}s_{\pi(1)},$$  

since each generator $s_i$ is an involution. Thus $\kappa(Y)$ is an upper bound for the number of conjugacy classes of Coxeter elements of a Coxeter group that has (unlabeled) Coxeter graph $Y$. This bound is known to be sharp in certain cases [15], but sharpness in the general case is still an open question.

Click-sequences are closely related to $c$-admissible sequences of a Coxeter element $c$. Recently, the structure of these sequences was studied and used to prove that a power of a Coxeter element of an infinite group is reduced [16].

The chip-firing game was introduced by Björner, Lovász, and Shor [2]. It is played over an undirected graph $Y$, and each vertex is given some number (possibly zero) of chips. If vertex $i$ has degree $d_i$, and at least $d_i$ chips, then a legal move (or a “click”) of vertex $i$ is a transfer of one chip to each neighboring vertex. This may be viewed as a generalization of a source-to-sink move for acyclic orientations where the out-degree of a vertex plays the role of the chip count. The chip-firing game is closely related to the numbers game [1]. In the numbers game over a graph $Y$, the legal sequences of moves are in 1–1 correspondence with the reduced words of the Coxeter group with Coxeter graph $Y$. For an excellent summary and comparison of these games, see [4].
A quiver is a finite directed graph (loops and multiple edges are allowed), and appears primarily in the study of representation theory. A quiver $Q$ with a field $K$ gives rise to a path algebra $KQ$, and there is a natural correspondence between $KQ$-modules, and $K$-representations of $Q$. In fact, there is an equivalence between the categories of quiver representations, and modules over path algebras. A path algebra is finite-dimensional if and only if the quiver is acyclic, and the modules over finite-dimensional path algebras form a reflective subcategory. A reflection functor maps representations of a quiver $Q$ to representations of a quiver $Q'$, where $Q'$ differs from $Q$ by a source-to-sink operation $\text{[10]}$. We note that while the composition of $n$ source-to-sink operations (one for each vertex) maps a quiver back to itself, the corresponding composition of reflection functors is not the identity, but rather a Coxeter functor. In fact, the same result in $\text{[10]}$ about powers of Coxeter elements being reduced was proven previously using techniques from the representation theory of quivers $\text{[6]}$.

We conclude with a remark on the evaluation of the functions $\alpha(Y)$ and $\kappa(Y)$. They both satisfy recurrences under edge deletion and contraction, and may be computed through evaluations of the Tutte polynomial $T_Y$ of $Y$. It is well-known that $\alpha(Y) = T_Y(2,0)$, and in $\text{[9]}$, we showed that $\kappa(Y) = T_Y(1,0)$. Other quantities counted by $T_Y(1,0)$ include the number of acyclic orientations of $Y$ with a unique sink at a fixed vertex $\text{[3]}$, and the Möbius invariant of the intersection lattice of the graphic hyperplane arrangement of $Y$ $\text{[12]}$. Some of the results in this paper have a natural interpretation in the language of the Tutte polynomial. For example, Corollary $\text{[1]}$ tells us that a connected graph $Y$ is bipartite if and only if $T_Y(1,0)$ is odd. Proposition $\text{[4]}$ implies that $n \cdot T_Y(1,0) \leq T_Y(2,0)$. We also point out a previous study of the Tutte polynomial in the context of the chip-firing game $\text{[7]}$. We hope this paper will motivate further explorations of the connections between these topics, as also provide insights useful to some of the open problems in Coxeter theory, in particular the sharpness of the bound $\kappa(Y)$ for the enumeration of conjugacy classes of Coxeter elements.

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