The Structure of ABC-minimal Trees with Given Number of Leaves

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Abstract

The atom-bond connectivity (ABC) index is a degree-based molecular descriptor with diverse chemical applications. Recent work of Lin et al. [21] gave rise to a conjecture about the minimum possible ABC-index of trees with a fixed number \( t \) of leaves. We show that this conjecture is incorrect and we prove what the correct answer is. It is shown that the extremal tree \( T_t \) is unique for \( t \geq 1195 \), it has order \(|T_t| = t + \lfloor \frac{t}{10} \rfloor + 1\) (when \( t \mod 10 \) is between 0 and 4 or when it is 5, 6, or 7 and \( t \) is sufficiently large) or \(|T_t| = t + \lfloor \frac{t}{10} \rfloor + 2\) (when \( t \mod 10 \) is 8 or 9 or when it is 5, 6, or 7 and \( t \) is sufficiently small) and its ABC-index is \( \left( \frac{10}{11} \sqrt{10} + \frac{1}{110} \sqrt{11} \right) \cdot t + O(1) \).

1 Introduction

One of the most important topological indices used in Chemical Graph Theory is the Atom Bond Connectivity index, also known as the ABC-index. It was introduced by Estrada [11] with relation to the energy of formation of alkanes. It was extensively studied in the last few years, from the point of view of chemical graph theory [12, 25], and in general graphs [6]. Additional chemical applications of the ABC-index were discovered [5, 7, 18, 20, 26]; we also refer to [8–10] and the references therein.

The ABC-index can be defined for any graph \( G \). For \( v \in V(G) \), let \( d_v \) denote the degree of the vertex \( v \). For each edge \( uv \in E(G) \), we consider its ABC-contribution, which is the value

\[
    f(d_u, d_v) = \sqrt{\frac{d_u + d_v - 2}{d_u d_v}}.
\]

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Then the Atom-Bond Connectivity index (shortly ABC-index) of $G$ is defined as the sum of ABC-contributions of its edges:

$$ABC(G) = \sum_{uv \in E(G)} f(d_u, d_v).$$

In this paper we shall only consider the ABC-indices of trees.

It is of interest to determine the extremal values of the ABC-index. In particular, the question about which trees attain the minimum possible value of the ABC-index among all trees with the given number of vertices has been extensively studied \[1, 4, 8, 10, 15, 17, 22\]; see also a survey article \[16\]. The solution to this long standing open question has been announced recently in \[19\].

A similar question has been investigated in \[13, 14, 21, 23, 24\], with intention to figure out which trees with the given number $t$ of leaves (i.e. vertices of degree 1) have the minimum value of the ABC-index. In this paper, a tree $T$ is said to be $t$-minimal if $T$ has $t$ leaves and no other tree with the same number of leaves has smaller ABC-index. The problem of classifying $t$-minimal trees has been raised in \[13\] and \[24\] and was further explored in \[14, 21, 23\]. Magnant et al. \[24\] claimed that $t$-minimal trees are balanced double stars whenever $t \geq 19$. This speculation was refuted by Goubko et al. in \[14\], where $t$-minimal trees were found for all values of $t$ up to 53. Lin et al. \[21\] improved the knowledge about $t$-minimal trees and were able to determine $t$-minimal trees for all $t \leq 219$. By exploring the patterns that have shown up in their computer-aided calculations, the following conjecture was proposed.

**Conjecture 1.1** (Lin et al. \[21\]). For $t \geq 88$, a $t$-minimal tree has $t + \lfloor \frac{t}{10} \rfloor - 1$ vertices.

Behind this conjecture there was also the precise description how the $t$-minimal trees would look like. Unfortunately, the authors of \[21\] failed to realize that with $t$ growing, the observed patterns may still change. Here we disprove their conjecture and determine the $t$-minimal trees for every $t \geq 1195$. Our main result is the following.

**Theorem 1.2.** For every $t \geq 1195$, there is a unique $t$-minimal tree $T_t$ that is described in the caption of Figures 1 and 2. Let $r = t - 10\lfloor \frac{t}{10} \rfloor$. If $0 \leq r \leq 7$ and $t$ fits the values listed at Figure 1, then $T_t$ has $t + \lfloor \frac{t}{10} \rfloor + 1$ vertices and its ABC-index is equal to

$$\left(\sqrt{\frac{10}{11}} + \sqrt{\frac{1}{10}}\right)t + 9\sqrt{\frac{1}{11}} + \left(11\sqrt{\frac{11}{12}} + \sqrt{\frac{1}{12}} - \sqrt{110} + \frac{1}{10}\sqrt{11}\right)r + O(t^{-1}).$$
If $5 \leq r \leq 9$ and $t$ fits the values listed at Figure 2, it has $t + \lfloor \frac{r}{10} \rfloor + 2$ vertices and its ABC-index is equal to

$$\left(\sqrt{\frac{10}{11}} + \frac{1}{10} \sqrt{\frac{1}{11}}\right)t + \frac{9}{2} \sqrt{\frac{1}{11}} + 9(10 - r)\left(\frac{14}{45} \sqrt{10} - \sqrt{\frac{10}{11}} - \frac{1}{10} \sqrt{\frac{1}{11}}\right) + O(t^{-1}).$$

**Figure 1.** The unique $t$-minimal tree with $t = 10p + r$ ($0 \leq r \leq 7$) has one root vertex, $p - r$ $S_{10}$-branches and $r$ $S_{11}$-branches (see Section 2 for definitions) when the following holds: ($r = 0$ and $t \geq 1030$) or ($r \in \{1, 2, 3, 4\}$ and $t \geq 1201$) or ($r = 5$ and $t \geq 1355$) or ($r = 6$ and $t \geq 2316$) or ($r = 7$ and $t \geq 7227$).

**Figure 2.** The unique $t$-minimal tree with $t = 10p + r$ ($5 \leq r \leq 9$) has one root vertex, $p + r - 9$ $S_{10}$-branches and $10 - r$ $S_{9}$-branches when ($r = 5$ and $1155 \leq t \leq 1345$) or ($r = 6$ and $1106 \leq t \leq 2306$) or ($r = 7$ and $1077 \leq t \leq 7217$) or ($r = 8$ and $t \geq 1058$) or ($r = 9$ and $t \geq 1039$).

In fact, the overall bound $t \geq 1195$ in Theorem 1.2 is best possible because the $t$-minimal tree for $t = 1194$ is a bit different. Our computer calculations show that it is soon after $t = 1000$ that we obtain structures as given in the theorem (see the caption of
the figures for a detailed information), but it is only after 1195 when we obtain this for all values of $t$ modulo 10.

The “evolution” of the structure of $t$-minimal trees is as follows:

(i) For $t \leq 18$, $t$-minimal trees are stars.

(ii) For $19 \leq t \leq 35$, $t$-minimal trees are balanced double stars, see [14].

(iii) From Goubko et al. [14], Lin et al. [21] and the calculations of this paper it follows that for $t \geq 36$, stars and double stars no longer occur. However, $t$-minimal trees have a mixed vertex for values of $t$ between 36 and 1194 (see definition of mixed vertices in the next section), except for those values that are indicated in captions of Figures 1 and 2.

(iv) For all values $t \geq 1195$, $t$-minimal trees are described by our Theorem 1.2.

2 Proof of Theorem 1.2

In our main proof, we will use the following results that are either well-known or easy to prove.

Lemma 2.1. Let $k$ be a fixed positive integer. The function $f_k(d) = f(d,k)$ has the following properties:

1. $f(d,1) = \sqrt{1 - \frac{1}{d}}$ is increasing in $d$ and $1 - \frac{1}{d} < f(d,1) < 1 - \frac{1}{2d}$.
2. $f(d,2) = \sqrt{\frac{1}{2}}$ is independent of $d$.
3. $f(d,k) = \sqrt{\frac{1}{k} + (1 - \frac{2}{k})\frac{1}{d}}$ is decreasing in $d$ for every fixed $k \geq 3$ and
   \[
   \sqrt{\frac{1}{k} (1 + \frac{k-2}{2d})} < f(d,k) < \sqrt{\frac{1}{k} (1 + \frac{k-2}{d})}.
   \]

Next we will discuss basic properties of $t$-minimal trees.

Lemma 2.2 ([21]). No $t$-minimal tree has a vertex of degree 2.

Lemma 2.3 ([19]). Let $uv$ and $u'v'$ be edges of a tree $T$. Let $T_v$ ($T_{v'}$) be the component of $T - uv$ ($T - u'v'$) containing the vertex $v$ ($v'$). Suppose that $T_v \cap T_{v'} = \emptyset$, and let $T'$ be the tree obtained from $T$ by exchanging the subtrees $T_v$ and $T_{v'}$. 

(a) If \( d_u > d_{u'} \) and \( d_v < d_{v'} \), then \( ABC(T) > ABC(T') \). In particular, \( T \) is not minimal.

(b) If \( d_u = d_{u'} \) or \( d_v = d_{v'} \), then \( ABC(T) = ABC(T') \).

Part (b) of the lemma motivates the following definition. We say that trees \( T \) and \( T' \) are similar (or ABC-similar) if \( T' \) can be obtained from \( T \) by a series of exchange operations, each of which satisfies one or the other equality of degrees in Lemma 2.3(b).

Note that similarity is an equivalence relation that preserves the ABC-index. In order to characterize \( t \)-minimal trees, it suffices to describe one element in each similarity class.

We will classify vertices of a tree into the following types:

(L) A vertex of degree 1 is a leaf.

(R) A vertex is a root if it is not adjacent to any leaf.

(S) A vertex of degree \( d > 1 \) is a star vertex if it is adjacent to \( k \geq d - 1 \) leaves\(^1\). A star vertex together with all adjacent leaves is a subtree and is said to be an \( S_k \)-branch or an \( S \)-branch of order \( k \).

(M) A vertex is a mixed vertex if it is adjacent to at least one leaf and to at least two non-leaf vertices.

**Lemma 2.4.** If \( T \) is a \( t \)-minimal tree, then it is similar to a tree with at most one mixed vertex.

**Proof.** Suppose that \( T \) is \( t \)-minimal and that it has two mixed vertices, \( u \) and \( u' \). If \( u \) and \( u' \) have different degrees, then we can make an exchange of a leaf at one of them with a larger degree subtree at the other vertex as in Lemma 2.3(a) and obtain a contradiction to minimality of \( T \). If \( d_u = d_{u'} \), then we can exchange leaves at one of these vertices with non-leaf subtrees (using the similarity exchange) and make one of \( u \) and \( u' \) either a star or a root vertex. By repeating this process, we arrive at a similar tree with at most one mixed vertex. \( \square \)

In the proofs below, we will compare the ABC-index of a tree \( T \) with that of a modified tree \( T' \). To make the notation shorter we will write

\[
\Delta(T, T') = ABC(T) - ABC(T').
\]

\(^1\)Note that \( k = d - 1 \) unless the tree is a star.
Lemma 2.5. If \( T \) is a \( t \)-minimal tree, then it has at most one root.

Proof. Let us first prove that no two root vertices can be adjacent. Suppose this is not the case and that \( x, y \) are two adjacent roots. Let \( x_1, \ldots, x_{d_x-1} \) be the neighbors of \( x \) different from \( y \) and let \( y_1, \ldots, y_{d_y-1} \) be the neighbors of \( y \) different from \( x \). Let \( T' \) be the tree obtained from \( T \) by contracting the edge \( xy \) and let \( w \) be the contracted vertex.

Then

\[
\Delta(T, T') = f(d_x, d_y) + \sum_{i=1}^{d_x-1} (f(d_x, d_{x_i}) - f(d_w, d_{x_i})) + \sum_{j=1}^{d_y-1} (f(d_y, d_{y_j}) - f(d_w, d_{y_j})).
\]

Each of the terms is positive by Lemma 2.1(3). This yields a contradiction to minimality of \( T \).

Thus, if there are two roots \( x, y \) in \( T \), there must be a mixed vertex \( z \) and both roots are adjacent to \( z \). We may assume that \( d_x \geq d_y \). Let \( z_0 \) be a leaf adjacent to \( z \). By using Lemma 2.3(a) we conclude that \( d_z \leq d_y \) since \( z \) has a neighbor of degree 1. Also, if \( d_z = d_y \), we can perform similarity exchanges at \( z \) and \( y \) so that \( z \) becomes a root vertex and \( y \) becomes a mixed vertex. However, this yields a contradiction since we would obtain a \( t \)-minimal tree with two adjacent roots. We conclude that \( d_x \geq d_y > d_z \).

Let us now consider the degree of \( x_1 \). Since \( d_x > d_z \), every neighbor of \( z \) different from \( x \) has degree at most \( d_{x_1} \) by Lemma 2.3(a). In particular, \( d_y \leq d_{x_1} \). Thus, \( d_z < d_{x_1} \). However, this gives a contradiction by Lemma 2.3(a) since \( x_1 \) is adjacent to a vertex of degree 1 and \( z \) is adjacent to \( y \) whose degree is more than 1. This contradiction completes the proof.

\[ \square \]

Lemma 2.6. Let \( T \) be a \( t \)-minimal tree. If a vertex \( v \) has as neighbors an \( S_k \)-branch and an \( S_l \)-branch, then \( |k - l| \leq 1 \).

Proof. Suppose that \( k \geq l + 2 \). Let \( d = d_v \). Lemma 2.3(a) implies that \( d \geq k + 1 \). By detaching a leaf from \( S_k \) and attaching it to \( S_l \), we obtain a tree \( T' \) in which \( S_k \) is replaced by \( S_{k-1} \) and \( S_l \) with \( S_{l+1} \). We have:

\[
\Delta(T, T') = f(d, k + 1) + f(d, l + 1) - f(d, k) - f(d, l + 2) + k f(k + 1, 1) + l f(l + 1, 1) - (k - 1) f(k, 1) - (l + 1) f(l + 2, 1).
\]

For any fixed \( k \) and \( l < k - 1 \), the value of the first four terms on the right-hand side of the above equality is increasing in terms of \( d \), so it suffices to treat the case when \( d \) is
minimum possible, \( d = k + 1 \). Then, for any fixed \( k \), the value is decreasing in \( l \), so we may assume that \( l = k - 2 \). In other words,

\[
\Delta(T, T') \geq f(k+1, k+1) + f(k+1, k-1) - 2f(k+1, k) + \nonumber
\]

\[
kf(k+1, 1) + (k-2)f(k-1, 1) - 2(k-1)f(k, 1).
\]

The right-hand side is decreasing in terms of \( k \) and is 0 in the limit when \( k \to \infty \). This shows that \( \Delta(T, T') > 0 \), which is a contradiction to the \( t \)-minimality of \( T \).

When \( t \leq 18 \), \( t \)-optimal trees are stars. For \( 19 \leq t \leq 35 \) then they become balanced double stars, see [14]. It follows from Goubko et al. [14] and Lin et al. [21] that for \( t \geq 36 \), stars and double stars no longer occur. Therefore, there exists a root or a mixed vertex. As a corollary, our lemmas above imply the following.

**Corollary 2.7.** If \( T \) is a \( t \)-minimal tree and \( t \geq 36 \), then \( T \) has either one mixed vertex, one root, or one mixed vertex and one root that are adjacent, and all other vertices are stars and leaves.

The results proved above give a restricted structure for \( t \)-minimal trees. The structure with both—a root and a mixed vertex—is shown in Figure 3. There may be just one root vertex (which turns out to be the case when \( t \geq 1195 \)) or just one mixed vertex (which happens when \( t < 1195 \) with several exceptions). The \( S \)-branches adjacent to \( R \) or to \( M \) have the same order or two consecutive orders (see Lemma 2.6). So we have up to 7 parameters: degrees \( d_R \) and \( d_M \) of \( R \) and \( M \), the number of leaves \( l \) at the mixed vertex,
the larger order $k_R$ and $k_M$ of $S$-branches at $R$ and at $M$, respectively, and the number $s_R$ and $s_M$ of $S$-branches of order $k_R - 1$ and $k_M - 1$. Clearly,
\[
t = l + (d_R - 1)k_R - s_R + (d_M - l - 1)k_M - s_M.
\]
In addition to this, we may assume (after using similarity exchanges if needed) to have the following inequalities (assuming $R$ and $M$ both exist):
\[
\begin{align*}
d_R &\geq d_M \geq k_R + 1 \geq k_M + 1 \\
k_R &\geq k_M \\
l &\leq d_M - 2 \\
s_R &\leq d_R - 2 \\
s_M &\leq d_M - l - 2.
\end{align*}
\]
Moreover, as proved in [21], $k_M \geq 5$. This enables us to make a brute force search for optimal parameters for every fixed $t$\footnote{When $t \leq 2000$, this takes only a couple of seconds on a desktop PC.}. Additional restrictions provided below reduce the number of cases to be treated and also establish our main theorem for large enough $t$. We will use the above notation in the rest of the paper. We shall also assume that $t \geq 36$.

Before continuing, we define the notion of the $ABC$-contribution $c(v)$ of a leaf $v$ in a tree $T$ as follows. If $v$ is contained in an $S_k$-branch and the star vertex is adjacent to a root or to a mixed vertex of degree $d$, then $c(v) = f(k+1, 1) + \frac{1}{k}f(k+1, d)$. If $v$ is a leaf adjacent to $M$ and $R$ exists, then $c(v) = f(d_M, 1) + \frac{1}{l}f(d_M, d_R)$. The remaining possibility is that $v$ is a leaf adjacent to $M$ and $R$ does not exists; in that case $c(v) = f(d_M, 1)$. Clearly, the sum of all contributions of the leaves satisfies:
\[
\sum_{v : \deg(v) = 1} c(v) = ABC(T). \tag{1}
\]

The leaf contributions of leaves adjacent to a star of order $k$ only depend on $k$ and on the degree $d$ of the root or mixed vertex adjacent to the star. In that case we also use the notation
\[
c(k, d) = f(k + 1, 1) + \frac{1}{k}f(k + 1, d).
\]
We will need the following basic properties of leaf contributions.

**Lemma 2.8.** The leaf contributions function $c(k, d)$ has the following properties.
(a) $c(k, d)$ is monotone decreasing in $d$ for every fixed $k \geq 2$.

(b) $c(k, d) - c(k_0, d)$ is monotone increasing in $d$ when $k > k_0$ and decreasing when $k < k_0$.

(c) The function $\Delta_0(k, d) = k \cdot c(k, d) - (k + 1)c(k + 1, d)$ is increasing in $d$.

Proof. (a) follows easily from Lemma 2.1(3). To prove (b) and (c), one simply looks at the derivatives

$$\frac{\partial}{\partial d} c(k, d) = \frac{1}{k} \frac{\partial}{\partial d} f(k + 1, d) = \frac{-\left(1 - \frac{2}{k+1}\right)}{2kd^2 f(k + 1, d)}$$

and

$$\frac{\partial}{\partial d} (k c(k, d)) = \frac{-\left(1 - \frac{2}{k+1}\right)}{2d^2 f(k + 1, d)}.$$

For (b) it suffices to consider the case when $k_0 = k - 1$ since $c(k, d) - c(k_0, d)$ can be written as the sum of consecutive differences of the form $c(i, d) - c(i - 1, d)$ for $i = k_0 + 1, \ldots, k$.

As this is a simple exercise, we leave the details to the reader. 

Table 1 shows contributions $c(k, d)$ of leaves in $S_k$-branches of order $k$ for $k = 5, \ldots, 16$ when the degree $d$ of the adjacent root or mixed vertex is 120 or very large (respectively). The differences between the values $c(k, d) - c(10, d)$ change when $d$ gets larger, but they stay between the two values in the table by Lemma 2.8(b). The minimum of $c(k, d)$ when $d \geq 120$ is fixed is always attained at $k = 10$.

| $k$ | $c(k, 120)$ | $c(k, 120) - c(10, 120)$ | $c(k, \infty)$ | $c(k, \infty) - c(10, \infty)$ |
|-----|-------------|--------------------------|-----------------|-------------------------------|
| 5   | 0.99587026  | 0.011146309              | 0.99452072      | 0.010906887                   |
| 6   | 0.99011316  | 0.005389211              | 0.9881431       | 0.005200473                   |
| 7   | 0.98716926  | 0.002445313              | 0.98592210      | 0.002308263                   |
| 8   | 0.98567376  | 0.000949811              | 0.98447583      | 0.000861993                   |
| 9   | 0.98497203  | 0.000248084              | 0.98381983      | 0.000205997                   |
| 10  | 0.98472395  | 0             | 0.9836138       | 0                             |
| 11  | 0.98474189  | 0.000017939            | 0.9836704       | 0.000056574                   |
| 12  | 0.98491753  | 0.000193580            | 0.9838815       | 0.000267700                   |
| 13  | 0.98518611  | 0.000462157            | 0.9841828       | 0.000568935                   |
| 14  | 0.98550786  | 0.000783911            | 0.9845347       | 0.000920824                   |
| 15  | 0.98585791  | 0.001133961            | 0.9849126       | 0.001298763                   |
| 16  | 0.98622049  | 0.001496542            | 0.9853011       | 0.001687235                   |

Table 1. Leaf contributions at $S_k$-branches (with the last shown digit rounded).

Lemma 2.9. Let $T$ be a $t$-minimal tree, where $t \geq 200$.

(a) If $T$ contains a root vertex, then the root is adjacent with an $S_k$-branch, where $k \leq 10$. Consequently, any $S_q$-branch in $T$ has $q \leq 11$. 

(b) If $T$ contains a mixed vertex, then it does not have a root vertex.

Proof. (a) If $T$ contains a root and a mixed vertex, then by using Lemma 2.3(a) we see that $T$ is similar to a minimal tree with $d_R \geq d_M \geq k_R + 1 \geq k_M + 1$, which we assume henceforth. This implies that more than $t/2$ leaves are within $S$-branches that are adjacent to the root $R$. To verify the claim, it suffices to see that there exists an $S$-branch of order 10 or less. Suppose, for a contradiction, that this is not the case. Then the $S$-branches adjacent to $R$ have orders $k_1 \geq k_2 \geq \cdots \geq k_{d_R-1} \geq 11$. Since $\sum_i k_i > \frac{t}{2} \geq 100$, it is easy to see that there is an index $j$ such that $\sum_{i=1}^j k_i \geq 10(j + 1)$. Then we may change $T$ to a tree with the same number of leaves by replacing the $S_{k_i}$-branches ($i = 1, \ldots, j$) with $j + 1$ branches, one of order 10 and the others of orders $k'_i$, where $10 \leq k'_i \leq k_i$ for each $i = 1, \ldots, j$. It is an easy calculation to show that the leaf contributions of all the leaves in $S$-branches adjacent to $R$ decrease. The degree of $R$ also increases, so the contribution of the possible edge $RM$ also drops (by Lemma 2.1(3)), and all other contributions remain unchanged. Therefore, the ABC-index drops, which is a contradiction.

(b) Suppose that $T$ has a root $R$ and a mixed vertex $M$. We will reach a contradiction in two steps. In the first step we show that $l < 37$.

Suppose that $l \geq 37$. We change the tree $T$ into a tree $T'$ which has one root vertex by moving all stars adjacent to $M$ to be adjacent to the root vertex and replacing $l$ pendant edges at $M$ with stars of orders 9, 10, 11 and make them adjacent to $R$ as well. By applying formula (1) and Lemma 2.1, we see that

$$\Delta(T, T') > lf(d_M, 1) - l \max_{k \in \{9,10,11\}} c(k, d_R + 4) \geq l(f(l + 2, 1) - c(9, l + 4)).$$

It is easy to see that the factor $f(l + 2) - c(9, l + 4)$ is positive for every $l \geq 37$. This contradicts $t$-minimality of $T$ and proves that $l \leq 36$.

Knowing that $l \leq 36$, we continue as follows. First, recall that $d_R \geq d_M \geq l + 2$. Thus, there are $l$ stars $S_1, \ldots, S_l$ adjacent to the root vertex. Each of them is of order $\leq 11$ by part (a). Now we change $T$ into a tree $T'$ as follows. We first move all stars adjacent to $M$ to the root vertex (which will be denoted by $R'$) and then, for $i = 1, \ldots, l$, replace each star $S_i$ of order $k_i$ with a star $S'_i$ of order $k_i + 1$. Then we remove the vertex $M$ and adjacent leaves. The resulting tree $T'$ has the same number of leaves and we will show that its ABC-index is smaller (thus giving a contradiction). The edge-contributions of leaves in all stars different from $S_1, \ldots, S_l$ have gone down (or stayed the same) because
\[ d_{R'} \geq d_R \text{ (see Lemma 2.8(a)). Thus} \]
\[
\Delta(T, T') \geq \sum_{i=1}^{l} \left( c(w_i) + k_i c(k_i, d_R) - (k_i + 1)c(k_i + 1, d_R') \right)
\]
where \( w_1, \ldots, w_l \) are the leaves adjacent to \( M \) in \( T \). We claim that each term in the above sum is positive.

First of all, we have:
\[
c(w_i) = f(d_M, 1) + \frac{1}{d_M}f(d_M, d_R)
\geq f(d_M, 1) + \frac{1}{d_M - 2}f(d_M, \infty)
= f(d_M, 1) + \frac{1}{d_M - 2}\sqrt{1/d_M}.
\]

By using Lemma 2.8(a) and (c) we conclude the following:
\[
k_i c(k_i, d_R) - (k_i + 1)c(k_i + 1, d_R') \geq k_i c(k_i, d_R') - (k_i + 1)c(k_i + 1, d_R')
\]
\[= \Delta_0(k_i, d_R') \geq \Delta_0(k_i, d_R).\]

By combining the above three inequalities, we see that it suffices to prove the following inequality for \( d = d_M \) and \( k = k_i \leq 11 \):
\[
f(d, 1) + \sqrt{\frac{1}{d(d-2)^2}} + \Delta_0(k, d) > 0. \tag{2}
\]

Inequality (2) has been verified for intermediate values of \( d \), \( 26 \leq d \leq 100 \), and for all \( k \) from 5 to 11 by a computer calculation. For \( d \geq 100 \) it can be proved analytically (the value is increasing in \( d \) for \( d \geq 100 \) for each \( k \)). The details are left to the reader.

Unfortunately, (2) fails for some values of \( k \leq 11 \) when \( d \leq 25 \). This range can be considered by using different approaches. We found it easiest to use our afore-mentioned algorithm to verify the claim by computer for all \( t \leq 1500 \). We may then assume that \( t \geq 1500 \). It is easy to see that in this case we have \( d_{R'} \geq 120 \). Moreover, by the same proof as used in part (a) of the proof, we see that there are more than \( l S_{10}^t \)-stars. Therefore, we only need to treat the case where \( k = 10 \). The inequality (2) can be replaced by:
\[
f(d, 1) + \sqrt{\frac{1}{d(d-2)^2}} + \Delta_0(10, 120) > 0. \tag{3}
\]
It is easy to see that (3) holds for every \( d \), \( 3 \leq d \leq 25 \). This completes the proof.

\[ \square \]

**Lemma 2.10.** If \( T \) is a \( t \)-minimal tree, where \( t \geq 1195 \), then \( T \) has a root vertex and does not have a mixed vertex, and \( T \) is isomorphic to the tree \( T_t \) as defined in Figures 1 and 2.
Proof. First, we claim that $T$ does not have a mixed vertex. If it does, there is no root vertex by the previous lemma. Let $M$ be the mixed vertex. Previous results imply that $d_M > 120$. Let $d'$ be the number of $S_{10}$-stars adjacent to $M$ and let $r$ be the number of $S_9$-stars or $S_{11}$-stars. Let $d = d_M - 1 = d' + r + l - 1 \geq d' + r$. The case when there are stars of order 9 is much easier to argue (by the same proof method as used below), thus we shall assume for brevity that we have no stars of order 9. It is easy to see that we cannot have ten or more $S_{11}$-stars (replacing 10 of them with 11 stars of order 10 decreases the ABC-index when $d_M > 120$). Thus $0 \leq r \leq 9$.

The following inequality which holds for every $d \geq 120$ is easy to verify:

$$f(11, d + 1) - f(11, d) + f(12, d + 1) - f(12, d) \leq 6 \cdot 10^{-4}. \quad (4)$$

Now we consider the following tree $T'$ with $t$ leaves. We remove one of the leaves adjacent to $M$ and change one $S_{10}$ into $S_{11}$. The following chain of inequalities use the following: Lemma 2.1(a) for the first inequality, (4) and $d' + r \leq d$ and $f(11, d + 1) - f(11, d) < 0$ and $r \leq 9$ for the second inequality:

$$\Delta(T, T') = 10f(1, 11) + f(1, d + 1) - 11f(1, 12) +$$
$$\quad (l - 1)(f(1, d + 1) - f(1, d)) +$$
$$\quad d'f(11, d + 1) - (d' - 1)f(11, d) +$$
$$\quad rf(12, d + 1) - (r + 1)f(12, d)$$
$$\geq 10f(1, 11) + f(1, d + 1) - 11f(1, 12) + f(11, d + 1) +$$
$$\quad (d' - 1)(f(11, d + 1) - f(11, d)) +$$
$$\quad r(f(12, d + 1) - f(12, d)) - f(12, d)$$
$$\geq 10f(1, 11) + f(1, d + 1) - 11f(1, 12) + f(11, d + 1) +$$
$$\quad (d - 1)(f(11, d + 1) - f(11, d)) - f(12, d) + 0.0054.$$ 

The last quantity above is positive for $d = 120$ and only increases when $d$ grows. This implies that $\Delta(T, T') > 0$. This contradiction shows that $T$ must have a root vertex and not a mixed vertex.

To deal with the case when there is a root vertex, we look at the contributions of the leaves. The minimum contribution is achieved with $S$-branches of order 10, the next

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3This conclusion needs some case analysis when $t$ is very close to 1195. Independently, we have checked our claims for small values of $t$ by computer and thus this is not really an issue to be worried about.

4This is easily tabulated for small values of $d \geq 120$. For the general case, basic calculus can be used to prove that the function is increasing.
smallest values are 11 and 9 (see Table 1), and by Lemma 2.6 we have only one of these. This means that the tree is one of those described in captions of Figures 1 and 2. This has been done by a computation for all values of \( t \leq 10000 \); the threshold values stated in Figures 1 and 2 have been obtained from these computations. For \( t \geq 10000 \) it suffices to prove that the following inequality (which compares the ABC-indices of trees from both figures) is satisfied for \( 0 \leq r \leq 7 \) and for \( p = \lfloor t/10 \rfloor \geq 1000 \):

\[
10(p - r)c(10, p) + 11r c(11, p) < 10(p + r - 9)c(10, p + 1) + 9(10 - r)c(9, p + 1)
\]

and that the reverse inequality holds when \( r = 8, 9 \). Again, this task reduces to verify the inequality for \( p = 1000 \) and then show that the difference is decreasing and that for \( r = 8, 9 \) the difference stays positive (by considering the limit when \( p \to \infty \)). It is easy to see that the worst cases are when \( r = 7 \) and \( r = 8 \) and that the proof for these two values implies the proof for all other values of \( r \).

**Lemma 2.11.** Suppose that \( t = 10p + r \), where \( 0 \leq r \leq 9 \).

(a) Let \( T_t \) be the tree as defined by Figure 1. Then

\[
\begin{align*}
\mathcal{ABC}(T_t) &= \left( \sqrt{\frac{10}{11}} + \frac{1}{10} \sqrt{\frac{9}{11}} \right) t + \frac{9}{2} \sqrt{\frac{1}{11}} + \\
&\quad \left( 11 \sqrt{\frac{11}{12}} + \sqrt{\frac{1}{12}} - \sqrt{\frac{110}{11}} + \frac{1}{10} \sqrt{\frac{1}{11}} \right) r + O(t^{-1}).
\end{align*}
\]

(b) Let \( T_t \) be the tree as defined by Figure 2. Then

\[
\begin{align*}
\mathcal{ABC}(T_t) &= \left( \sqrt{\frac{10}{11}} + \frac{1}{10} \sqrt{\frac{9}{11}} \right) t + \frac{9}{2} \sqrt{\frac{1}{11}} + \\
&\quad 9(10 - r) \left( \frac{14}{35} \sqrt{10} - \sqrt{\frac{10}{11}} - \frac{1}{10} \sqrt{\frac{1}{11}} \right) + O(t^{-1}).
\end{align*}
\]

*Proof.* The formula (1) implies, for the first case, that

\[
\mathcal{ABC}(T_t) = 10(p - r)c(10, p) + 11r c(11, p)
\]

and for the second one:

\[
\mathcal{ABC}(T_t) = 10(p + r - 9)c(10, p + 1) + 9(10 - r)c(9, p + 1).
\]

A routine calculation using approximations

\[
f(k, d) = \sqrt{\frac{1}{k}} \left( 1 + \frac{k - 1}{2d} - O(d^{-2}) \right) \quad \text{and} \quad d = \frac{1}{10} + O(1)
\]

(see Lemma 2.13) gives the claimed expressions.  \(\square\)
Proof of Theorem 1.2. The structure of $t$-minimal trees described after Corollary 2.7 enabled us to search for minimal trees for all values of $t \leq 2000$. The calculations verify the claim of the theorem. On the other hand, for $t \geq 2000$, the results given above show the same: we have one root vertex and any $t$-minimal tree is isomorphic to $T_t$. Finally, Lemma 2.11 gives the asymptotic value of $ABC(T_t)$.

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