KEN KUNEN: ALGEBRAIST

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1. Introduction

Ken Kunen is justifiably best known for his work in set theory and topology. What I would guess many of his friends and students in those areas do not know is that Ken also did important work in algebra, especially in quasigroup and loop theory. In fact, I think it is not an exaggeration to say that his work in loop theory, both alone and in collaboration, revolutionized the field. In this paper, I would like to describe some of his accomplishments to nonspecialists. My point of view is personal, of course, and so I will give the most attention to those of his projects in which he collaborated with me. My hope is that the set theorists and topologists reading this will come away with an appreciation for what Ken was able to do in an area outside of his direct speciality. In the interest of space, I will have to leave out discussion of some of Ken’s work, such as his paper on alternative loop rings [15].

Ken’s approach to algebra utilized automated deduction tools and finite model builders. At the time most of what I am going to describe took place, the automated deduction software of choice was OTTER, developed by William McCune [18]. (McCune is best known to mathematicians for his solution to the Robbins Problem in Boolean algebras [20].) The finite model builder Ken used during the period I will discuss was SEM, developed by J. Zhang and H. Zhang [23]. In recent years, these have been supplanted by other tools, such as McCune’s PROVER9, a successor to OTTER, and McCune’s model builder MACE4 [19].

2. Single Axioms for Groups

Of course, in my choices of which parts of Ken’s work to discuss here, I do not mean to suggest that his work in set theory and topology is entirely devoid of algebraic content. But I think it is reasonable to say that Ken’s purely algebraic work begins with his papers [11, 12] and also [10] written in collaboration with Joan Hart.

Ken’s work on single axioms for groups followed up earlier work of McCune and Wos (see the bibliographies of [11] [12] [10]). McCune had shown that

\[ w((x^{-1}w)^{-1}z)((yz)^{-1}y) = x \]

is a single axiom for group theory in the language of one binary operation \( \cdot \) (written as juxtaposition) and one unary operation \( ^{-1} \). In [11], Ken showed that McCune’s axiom is the shortest of its particular type, consisting of 7 variable occurrences on the left side and 4 variables. Ken also showed that

\[ (yy^{-1})^{-1}[(y^{-1}z)((yx)^{-1}z)^{-1}] = x \]

is a shortest single axiom for group theory of the type with 7 variable occurrences on the left and 3 variables. The follow-up papers [12] [10] gave single axioms for groups of exponent 4 and odd exponent, respectively.
Already in these papers we see Ken’s nascent interest in nonassociative structures. In order to verify that certain axioms are optimal, he found it necessary to construct nonassociative models showing that other possible candidate single axioms did not work. In retrospect, these models turned out to be loops. Ken told me later that at the time he was writing these papers, he did not even know what a loop was!

It is also in the single axioms papers that we find the beginnings of what I will call the *Humanization Imperative*:

*Proofs generated by automated deduction tools should be humanized, if it is reasonable to do so.*

Disclaimer: Ken never stated this imperative in precisely this form, nor even called it an imperative. Nevertheless, we find the idea in these early algebra papers. What follows is a bit of context for what the imperative means.

Most papers, both then and now, in which automated deduction tools are applied to mathematical problems do not necessarily give much human understanding of the proof the tools generated. Generally, the problem and the strategies for using the tools are described and the result is announced. Sometimes, the software’s proof will be included in the paper if it is not too long; otherwise, it will be left to a companion website or available from the author upon request. Of course, such proofs are usually not conceptual, particularly those involving convoluted equational reasoning. A human reading one of these papers is unlikely to come away with any understanding beyond a sense that a proof exists.

What Ken began to articulate in the single axiom papers is that it is incumbent upon those of us who use automated deduction tools to try to find conceptual, humanly understandable proofs, whenever doing so is reasonable. Here is an explicit statement of this taken from [10]:

> We...used OTTER as a reasoning assistant...but found that by examining the output from our assistant, we could provide conceptual proofs which a human could also understand. This conceptual understanding, in turn, led us to discover more axioms.

The Humanization Imperative plays a role in Ken’s work in quasigroup and loop theory as well.

### 3. Moufang Quasigroups

Combinatorially, a *quasigroup* \((Q, \cdot)\) is a set \(Q\) with a binary operation \(\cdot\) (represented by juxtaposition) such that for each \(a, b \in Q\), the equations \(ax = b\) and \(ya = b\) have unique solutions \(x, y \in Q\). In the finite case, the multiplication tables of quasigroups are Latin squares. A *loop* is a quasigroup with a neutral element: \(1x = x1 = x\) for each \(x\). A standard reference for quasigroup and loop theory is [2], and any uncited concepts in this paper can be found there.

A somewhat more useful definition is that preferred by universal algebraists. To them, a quasigroup has not one, but three binary operations \(\cdot, \backslash, /\) satisfying the identities

\[
\begin{align*}
x\backslash(xy) &= y \\
(xy)/y &= x \\
x(x\backslash y) &= y \\
(x/y)y &= x.
\end{align*}
\]
A Moufang loop is one satisfying the Moufang identities
\[ x((yz)x) = (xy)(zx) \quad (x(yz))x = (xy)(zx) \]
\[ x(yxz) = ((xy)x)z \quad ((xy)z)y = x(yzy) \].

These identities turn out to be equivalent in loops; that is, if a loop satisfies any one of them, then it satisfies all four.

A classical result in quasigroup theory is that an associative quasigroup is a group. This is equivalent to proving that an associative quasigroup has a neutral element, that is, it is a loop. Equationally, this can be expressed in quasigroups by the identity \( x_\cdot x = y/y \). The proof is easy: Start with \( xy = (x(x\backslash x))y = x((x\backslash x)y) \), then cancel \( x \)'s on the left to get \( y = (x\backslash x)y \), from which the desired identity follows immediately.

In [13], Ken extended this result by showing that a quasigroup satisfying any one of the Moufang identities is a loop, that is, must have a neutral element. It follows that the Moufang identities are not only equivalent in loops, they are, in fact, equivalent in quasigroups. In [14], he extended this to other classes of quasigroups and loops satisfying so-called identities of Bol-Moufang type. For example, extra loops are those satisfying the identities
\[ (x(yz))y = (xy)(zy) \quad ((xy)z)x = x(yzx) \quad (yz)(yx) = y((zy)x) \].

These identities are equivalent in loops, and once again, Ken showed they are equivalent in quasigroups.

4. Conjugacy closed loops and G-loops

For each \( x \) in a loop \( Q \), define the left and right translation maps \( L_x : Q \rightarrow Q \) and \( R_x : Q \rightarrow Q \) by \( yL_x = xy \) and \( yR_x = yx \). (Here we follow the convention that permutations act on the right of their arguments, the convention Ken followed in his papers.) A loop \( Q \) is said to be conjugacy closed (or a CC-loop, for short) if the sets \( \{ L_x \mid x \in Q \} \) and \( \{ R_x \mid x \in Q \} \) are each closed under conjugation. Thus for each \( x, y \in Q \), there exists \( z \in Q \) such that \( L^{-1}_x L_y L_x = L_z \) and similarly for the right translations. These conditions can be easily expressed purely equationally, thus making them ripe for exploration using automated deduction.

A (principal) isotope of a loop \( (Q, \cdot) \) is the same set \( Q \) with another loop operation \( \circ \) defined by \( x \circ y = (x/a)(b\backslash y) \) where \( a, b \) are fixed elements of \( Q \). Isotopes are of no interest to group theorists: two groups are isotopic if and only if they are isomorphic. However, this property does not characterize groups. There exist nonassociative loops which are isomorphic to all of their isotopes; such loops are known as G-loops. (The term is imported from the Russian literature; “G” is probably short for “grouplike”.) CC-loops, for instance, are G-loops, but there are other types as well.

CC-loops were introduced in the west by Goodaire and Robinson [4], but went somewhat unexplored for about 17 years. Ken revived interest in these loops in [17]. That paper gives a detailed structure theory (partially subsumed by later work), and also constructs examples of CC-loops, particularly “small” ones (e.g., of order \( pq \), etc.).

In that same paper, he also showed that the equational theory of CC-loops does not extend to all G-loops, because any identity holding in all G-loops holds in all loops. He did this by constructing a “universal” G-loop which contains isomorphic copies of all finite loops.
In the follow-up paper [16] (which actually appeared first, because of journal backlogs), Ken studied G-loops in more detail and showed that there are no nonassociative examples of order $3q$ where $q$ is prime and $3$ does not divide $q - 1$.

5. Diassociative A-loops

Now we turn to the more personal side of the story. First, a bit of mathematical background. The multiplication group $\text{Mlt}(Q)$ of a loop $Q$ is the group generated by all left and right translation maps. Its stabilizer of the neutral element is called the inner mapping group $\text{Inn}(Q)$. This is a natural generalization of the inner automorphism group of group theory. The inner mapping group is generated by all mappings of the forms $R_xL_y^{-1}$ (measures of noncommutativity, just as we would expect), $R_xR_yR_y^{-1}$ and $L_xL_yL_y^{-1}$ (measures of nonassociativity).

Unlike the group case, inner mappings of a loop need not act as automorphisms of the loop. A loop for which this does occur is called an A-loop, a notion introduced by Bruck and Paige in 1956 [3]. (These are now sometimes called automorphic loops to avoid the excessive acronymization that has plagued the field.) Every commutative Moufang loop is an A-loop, for instance.

A loop $Q$ is diassociative if given any $x, y \in Q$, the subloop generated by $x$ and $y$ is a group. Every Moufang loop is diassociative; this is, in fact, a weak form of what is usually known as Moufang’s theorem.

Bruck and Paige studied diassociative A-loops, and noted that they seem to have many Moufang-like properties. They did not explicitly state a conjecture, but the subtext of the paper makes it clear that they were conjecturing that every diassociative A-loop is a Moufang loop. J. M. Osborn showed that this conjecture is true for commutative A-loops [21]. Osborn’s paper is a triumph of human equational reasoning.

The general conjecture sat dormant for many years, and it is probably fair to say that most loop theorists forgot about A-loops. Someone who was aware of it was Tomáš Kepka of Charles University in Prague. He suggested it as an interesting problem to J. D. Phillips (then of St. Mary’s College in Moraga, now at Northern Michigan University).

In 1999, J. D. and Hala Pflugfelder and I were having an email conversation about every loop-theoretic topic under the sun. J. D. told me about the Bruck-Paige problem, and the two of us discussed it at length, cc’ing Hala each time. Finally, Hala suggested that we should ask Ken Kunen. J. D.’s reaction was “the set theorist?” I had not had a set theory class when I was a graduate student, so my reaction was “who?” Hala explained to us that Ken had been getting some very interesting results using computers.

Somewhat brashly, I sent an email message to Ken describing the problem to him. Because of the explicit generators of $\text{Inn}(Q)$ mentioned above, the problem has an explicit equational form. Thus it can be formulated as a problem for automated deduction tools. Ken’s email reaction was classic: Interesting… . I later learned that “Interesting…” (with the ellipsis) was a clue that I had hooked him with a problem.

Ken set up the problem for OTTER and started working on it over the course of a few days. (At the time, neither J. D. nor I knew how to use OTTER.) About four or five days after I originally posed the problem to him, J. D. and I received an email message with roughly the following header:
And so it was. OTTER had confirmed the Bruck-Paige conjecture, and the proof was indeed quite ugly.

By this time, I had caught up with reading Ken’s earlier loop theory papers, and was familiar with what I have been calling here the Humanization Imperative. So when J. D. asked me privately what the protocol was, I told him that now our job was to translate the proof into human form.

Ken and I actually made the mistake of starting at the front of the proof and working forward. J. D. started at the end of the proof and worked backward, eventually reaching a point where the proof relied on results which were already in the Bruck-Paige paper. After a couple of rounds of polishing, the paper was ready. Although most of the argument is motivated, at the very center of the proof is an interesting piece of equational reasoning which, to me, at least, does not seem to be something which could have been discovered by humans. It is nevertheless easy for a human to verify, unlike the original OTTER output.

We submitted the paper [5] to the Proceedings of the AMS, the same journal that had published Osborn’s paper some 41 years earlier. It was accepted for publication five days after we submitted it!

6. The KK(P) Collaboration

From that point on, Ken’s work in loop theory was joint with me and J. D. or just me. In [6], we extended Moufang’s theorem to a very wide class of loops which include both Moufang loops and Steiner loops (which arise from Steiner triple systems) as special cases. In this case, we did not obtain the main result directly from OTTER, but we used OTTER to verify enough instances of the result that we were able to construct an inductive argument.

In [7], we returned to CC-loops. Quite by accident, I discovered that before Goodaire and Robinson introduced them to the western world, they were already known in the Russian literature thanks to work of Soikis [22]. In particular, a key open question in the Goodaire-Robinson paper, restated in Ken’s CC-loops paper, was solved in 1991 by A. Basarab [1]: the factor loop of a CC-loop by its nucleus is an abelian group. (Informally, the nucleus of a loop is the set of all elements that associate with all elements of the loop.) It is interesting to speculate how Ken’s own CC-loop paper would have changed had he known about Basarab’s earlier result. In any case, we used Basarab’s result along with Ken’s structure theory to study particular instances of diassociativity inside of CC-loops.

CC-loops which are fully diassociative are precisely the extra loops studied earlier. Ken and I worked on these before, during and after we met in person at the LOOPS conference in Prague in 2003. We developed a very detailed structure theory for them [8]. Later, we worked out a similar theory for the more general class of power-associative CC-loops [9]. (A loop is power-associative if each 1-generated subloop is a group.) Even now, loop theorists have told me how surprising they find the main result of that paper to be: the factor loop of a power-associative, CC-loop by its nucleus is an abelian group of exponent 12.
7. Epilogue

After [9], Ken mostly went back to his set-theoretic and topological roots. (We have a couple of other projects languishing on our back burners.) In the meantime, both J. D. and I had started using OTTER and its successor Prover9 in our own loop theory work. And it is actually this for which I am the most grateful to Ken. He was very patient with all of my stupid questions about why certain input files were wrong, what various output files meant, and so on.

I hope I have convinced the reader of the significance of Ken’s work in quasigroup and loop theory. He introduced a new methodological paradigm: not only can and should we use automated deduction tools to assist us, but we should also work to understand what those tools tell us. For this, the field owes Ken a great debt.

I congratulate Ken on his change of employment status, and wish him all the best in the years to come.

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