Matrix Completion has No Spurious Local Minimum

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Abstract
Matrix completion is a basic machine learning problem that has wide applications, especially in collaborative filtering and recommender systems. Simple non-convex optimization algorithms are popular and effective in practice. Despite recent progress in proving various non-convex algorithms converge from a good initial point, it remains unclear why random or arbitrary initialization suffices in practice. We prove that the commonly used non-convex objective function for matrix completion has no spurious local minima – all local minima must also be global. Therefore many popular optimization algorithms such as (stochastic) gradient descent can provably solve matrix completion with arbitrary initialization in polynomial time.

1 Introduction
Matrix completion is the problem of recovering a low rank matrix from partially observed entries. It has been widely used in collaborative filtering and recommender systems [Kor09, RS05], dimension reduction [CLMW11] and multi-class learning [AFSU07]. There has been extensive work on designing efficient algorithms for matrix completion with guarantees. One earlier line of results (see [Rec11, CT10, CR09] and the references therein) rely on convex relaxations. These algorithms achieve strong statistical guarantees, but are quite computationally expensive in practice.

More recently, there has been growing interest in analyzing non-convex algorithms for matrix completion [KMO10, JNS13, Har14, HW14, SL15, ZWL15, CW15]. Let $M \in \mathbb{R}^{d \times d}$ be the target matrix with rank $r \ll d$ that we aim to recover, and let $\Omega = \{(i, j) : M_{i,j} \text{ is observed}\}$ be the set of observed entries. These methods are instantiations of optimization algorithms applied to the objective

$$f(X) = \frac{1}{2} \sum_{(i,j) \in \Omega} \left[M_{i,j} - (XX^\top)_{i,j}\right]^2,$$

These algorithms are much faster than the convex relaxation algorithms, which is crucial for their empirical success in large-scale collaborative filtering applications [Kor09].

All of the theoretical analysis for the nonconvex procedures require careful initialization schemes: the initial point should already be close to optimum. In fact, Sun and Luo [SL15] showed that after this initialization the problem is effectively strongly-convex, hence many different optimization procedures can be analyzed by standard techniques from convex optimization.

However, in practice people commonly use a random initialization, which still leads to robust and fast convergence. How can these practical algorithms find the optimal solution in spite of the non-convexity? In this work we investigate this question and show that the matrix completion objective has no spurious local minimum. More precisely, we show that any local minimum $X$ of objective function $f(\cdot)$ is also a global minimum with $f(X) = 0$, and recovers the correct low rank matrix $M$.

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1In this paper we focus on the symmetric case when true $M$ has symmetric decomposition $M = ZZ^\top$. Some of previous papers work on the asymmetric case when $M = ZW^\top$, which is relatively harder than the symmetric case.

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Our characterization of the structure in the objective function implies that (stochastic) gradient descent from arbitrary starting point converge to a global minimum. This is because gradient descent converges to a local minimum \cite{GHJY15, LSJR16}, and for our objective function a local minimum is also a global one.

1.1 Main results

Assume the target matrix $M$ is symmetric and each entry of $M$ is observed with probability $p$ independently\footnote{The entries $(i, j)$ and $(j, i)$ are the same. With probability $p$ we observe both entries and otherwise we observe neither.}. We assume $M = ZZ^T$ for some matrix $Z \in \mathbb{R}^d$.

There are two known issues with matrix completion. First, the choice of $Z$ is not unique since for any orthonormal matrix $R \in \mathbb{R}^{d \times r}$, we have $M = (ZR)(ZR)^\top$. Our goal is to find one of these equivalent solutions.

Another issue is that matrix completion is impossible when $M$ is “aligned” with standard basis. For example, when $M$ is the identity matrix in its first $r \times r$ block, we will very likely be observing only 0 entries. To address this issue, we make the following standard assumption:

**Assumption 1.** For any row $Z_i$ of $Z$, we have

$$\|Z_i\| \leq \frac{\mu}{\sqrt{d}} \cdot \|Z\|_F.$$  

Moreover, $Z$ has a bounded condition number $\sigma_{\max}(Z)/\sigma_{\min}(Z) = \kappa$.

Throughout this paper we think of $\mu$ and $\kappa$ as small constants, and the sample complexity depends polynomially on these two parameters. Also note that this assumption is independent of the choice of $Z$: all $Z$ such that $ZZ^T = M$ have the same row norms and Frobenius norm.

This assumption is similar to the “incoherence” assumption \cite{CR09}. Our assumption is the same to the ones used in analyzing non-convex algorithms \cite{KMO10, SL15}.

We enforce $X$ to also satisfy this assumption by a regularizer

$$f(X) = \frac{1}{2} \sum_{i,j \in \Omega} \left[ M_{i,j} - (XX^T)_{i,j} \right]^2 + R(X),$$  \hspace{1cm} (1.2)

where $R(X)$ is a function that penalizes $X$ when one of its rows is too large. See Section 3 for the precise definition. Our main result shows that in this setting, the regularized objective function has no spurious local minimum:

**Theorem 1.1.** [Informal] All local minimum of the regularized objective (1.1) satisfy $XX^T = ZZ^T = M$ when $p \geq \text{poly}(\kappa, r, \mu, \log d)/d$.

Combined with the results in \cite{GHJY15, LSJR16} (see more discussions in Section 1.2), we have,

**Theorem 1.2** (Informal). With high probability, stochastic gradient descent on the regularized objective (1.1) will converge to a solution $X$ such that $XX^T = ZZ^T = M$ in polynomial time from any starting point. Gradient descent will converge to such a point with probability 1 from a random starting point.

Our results are also robust to noise. Even if each entry is corrupted with Gaussian noise of standard deviation $\mu^2 \|Z\|_F^2 / d$ (comparable to the entry itself!), we can still guarantee that all the local minima satisfy $\|XX^T - ZZ^T\|_F \leq \epsilon$ when $p$ is large enough, see more discussions in Appendix B.

Our main technique is to show that every point that satisfies the first and second order necessary conditions for optimality must be a desired solution. To achieve this we use new ideas to analyze the effect of the regularizer and show how it is useful in modifying the first and second order conditions to exclude any spurious local minimum.
1.2 Related Work

Matrix Completion. The earlier theoretical works on matrix completion analyzed the nuclear norm heuristic [Rec11, CT10, CR09]. This line of work has the cleanest and strongest theoretical guarantees; [CT10, Rec11] showed that if \(|\Omega| \geq d^{2}r^{2}\log^{2}d\) the nuclear norm convex relaxation recovers the exact underlying low rank matrix. The solution can be computed via the solving a convex program in polynomial time. However the primary disadvantage of nuclear norm methods is their computational and memory requirements. The fastest known algorithms have running time \(O(d^{3})\) and require \(O(d^{2})\) memory, which are both prohibitive for moderate to large values of \(d\). These concerns led to the development of the low-rank factorization paradigm of [BM03]: Burer and Monteiro proposed factorizing the optimization variable \(\tilde{M} = XX^T\), and optimizing over \(X \in \mathbb{R}^{d \times r}\), instead of \(M \in \mathbb{R}^{d \times d}\). This approach only requires \(O(dr)\) memory, and a single gradient iteration takes time \(O(r|\Omega|)\), so has much lower memory requirement and computational complexity than the nuclear norm relaxation. On the other hand, the factorization causes the optimization problem to be non-convex in \(X\), which leads to theoretical difficulties in analyzing algorithms. Under incoherence and sufficient sample size assumptions, [KMO10] showed that well-initialized gradient descent recovers \(M\). Similarly, [HW14, Har14, INS13] showed that well-initialized alternating least squares or block coordinate descent converges to \(M\), and [CW15] showed that well-initialized gradient descent converges to \(M\). [SL15, ZWL15] provided a more unified analysis by showing that with careful initialization many algorithms, including gradient descent and alternating least squares, succeed. [SL15] accomplished this by showing an analog of strong convexity in the neighborhood of the solution \(M\).

Non-convex Optimization. Recently, a line of work analyzes non-convex optimization by separating the problem into two aspects: the geometric aspect which shows the function has no spurious local minimum and the algorithmic aspect which designs efficient algorithms can converge to local minimum that satisfy first and (relaxed versions) of second order necessary conditions.

Our result is the first that explains the geometry of the matrix completion objective. Similar geometric results are only known for a few problems: phase retrieval/synchronization, tensor decomposition, dictionary learning [GHJY15, SQW15, BBV16]. The matrix completion objective requires different tools due to the sampling of the observed entries, as well as carefully managing the regularizer to restrict the geometry. Parallel to our work Bhojanapalli et al. [BNS16] showed similar results for matrix sensing, which is closely related to matrix completion.

On the algorithmic side, it is known that second order algorithms like cubic regularization [NP00] and trust-region [SQW15] algorithms converge to local minima that approximately satisfy first and second order conditions. Gradient descent is also known to converge to local minima [LSJR16] from a random starting point. Stochastic gradient descent can converge to a local minimum in polynomial time from any starting point [Pem90, GHJY15]. All of these results can be applied to our setting, implying various heuristics people use in practice are guaranteed to solve matrix completion.

2 Preliminaries

2.1 Notations

For \(\Omega \subset [d] \times [d]\), let \(P_\Omega\) be the operator that maps a matrix \(A\) to \(P_\Omega(A)\), where \(P_\Omega(A)\) has the same values as \(A\) on \(\Omega\), and 0 outside of \(\Omega\).

We will use the following matrix norms: \(\|\cdot\|_F\) the frobenius norm, \(\|\cdot\|\) spectral norm, \(|A|_\infty\) elementwise infinity norm, and \(|A|_{p \to q} = \max_{|\alpha| = 1} \|A\|_q\). We use the shorthand \(\|A\|_\Omega = \|P_\Omega A\|_F\). The trace inner product of two matrices is \(\langle A, B \rangle = \text{tr}(A^T B)\), and \(\sigma_{\min}(X), \sigma_{\max}(X)\) are the smallest and largest singular values of \(X\). We also use \(X_i\) to denote the \(i\)-th row of a matrix \(X\).

2.2 Necessary conditions for Optimality

Given an objective function \(f(x) : \mathbb{R}^n \rightarrow \mathbb{R}\), we use \(\nabla f(x)\) to denote the gradient of the function, and \(\nabla^2 f(x)\) to denote the Hessian of the function \((\nabla^2 f(x))_{i,j} = \frac{\partial^2}{\partial x_i \partial x_j} f(x)\). It is well known that local minima of the function \(f(x)\) must satisfy some necessary conditions:
Definition 2.1. A point $x$ satisfies the first order necessary condition for optimality (later abbreviated as first order optimality condition) if $\nabla f(x) = 0$. A point $x$ satisfies the second order necessary condition for optimality (later abbreviated as second order optimality condition) if $\nabla^2 f(x) \succeq 0$.

These conditions are necessary for a local minimum because otherwise it is easy to find a direction where the function value decreases. We will also consider a relaxed second order necessary condition, where we only require the smallest eigenvalue of the Hessian $\nabla^2 f(x)$ to be not very negative:

Definition 2.2. For $\tau \geq 0$, a point $x$ satisfies the $\tau$-relaxed second order optimality condition, if $\nabla^2 f(x) \succeq -\tau \cdot I$.

This relaxation to the second order condition makes the conditions more robust, and allows for efficient algorithms.

Theorem 2.3. \cite{NP06, SQW15, GHJY15} If every point $x$ that satisfies first order and $\tau$-relaxed second order necessary condition is a global minimum, then many optimization algorithms (cubic regularization, trust-region, stochastic gradient descent) can find the global minimum up to $\epsilon$ error in function value in time $\text{poly}(1/\epsilon, 1/\tau, d)$.

3 Warm-up: Rank-1 matrix completion

In this section we analyze the geometry of the objective function for rank $r = 1$ case with a simple and clean proof. This case illustrates our main ideas. The rank $r$ analysis follows from the same approach and is shown in the next section.

In this case, assume $M = zz^T$, where $||z|| = 1$, and $||z||_\infty \leq \frac{d}{\sqrt{d}}$. The objective function simplifies to,

$$f(x) = \frac{1}{2} ||P_\Theta(M - xx^T)||^2_F + \lambda R(x).$$

(3.1)

Here we use the the regularization $R(x)$

$$R(x) = \sum_{i=1}^d h(x_i), \text{ and } h(t) = (|t| - \alpha)^4 1_{|t| \geq \alpha}.$$ 

The parameters $\lambda$ and $\alpha$ will be chosen later as in Theorem 3.2. We will choose $\alpha > 10\mu / \sqrt{d}$ so that $R(x) = 0$ for incoherent $x$, and thus it only penalizes coherent $x$. Moreover, we note that $R(x)$ has Lipschitz second order derivative.

We first state the optimality conditions, whose proof is deferred to Appendix A.

Proposition 3.1. The first order optimality condition of objective \eqref{3.1} is,

$$2P_\Theta(M - xx^T)x = \lambda \nabla R(x),$$

and the second order optimality condition requires:

$$\forall v \in \mathbb{R}^d, \ |P_\Theta(v x^T + xv^T)||^2_F + \lambda v^T \nabla^2 R(x) v \geq 2v^T P_\Theta(M - xx^T) v.$$

(3.3)

Moreover, The $\tau$-relaxed second order optimality condition requires

$$\forall v \in \mathbb{R}^d, \ |P_\Theta(v x^T + xv^T)||^2_F + \lambda v^T \nabla^2 R(x) v \geq 2v^T P_\Theta(M - xx^T) v - \tau ||v||^2.$$

(3.4)

We give the precise version of Theorem 1.1 for the rank-1 case.

Theorem 3.2. For $p \geq \frac{c^2 \log^{13/4} d}{d}$ where $c$ is a large enough absolute constant, set $\alpha = 10\mu \sqrt{1/\tau}$ and $\lambda \geq \mu^2 p/\alpha^2$. Then, with high probability over the randomness of $\Omega$, the only points in $\mathbb{R}^d$ that satisfy first and (or $\tau$-relaxed with $\tau < -0.1p$) second order optimality conditions are $z$ and $-z$. 

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If we observe every entry of $M$, Theorem 3.2 follows from simple linear algebra: the first order optimality condition would require $Mx = \|x\|^2 x$, and therefore $x$ has to be an eigenvector of $M$. The second order optimality condition will make sure $x$ is the eigenvector with largest eigenvalue. When we only observe a small subset of entries, we will use concentration bounds to prove more robust versions of these arguments. In the rest of this section, we will first prove that when $x$ is constrained to be incoherent (and hence the regularizer is 0 and concentration is straightforward) and satisfies the optimality conditions, then $x$ has to be $z$ or $-z$. Then we go on to explain how the regularizer helps us to change the geometry of those points that are far away from $z$ so that we can rule out them from being local minimum.

For simplicity, we will focus on the part that shows a local minimum $x$ must be close enough to $z$.

**Lemma 3.3.** In the setting of Theorem 3.2, suppose $x$ satisfies the first-order and second-order optimality condition (3.2) and (3.3). Then when $p$ is defined as in Theorem 3.2,

$$\|xx^T - zz^T\|_F \leq O(\epsilon),$$

where $\epsilon = \mu^3 (pd)^{-1/2}$.

This turns out to be the main challenge. Once we proved $x$ is close, we can apply the result of Sun and Luo [SL15] (see Lemma C.1), and obtain Theorem 3.2.

### 3.1 Handling incoherent $x$

To demonstrate the key idea, in this section we restrict our attention to the subset of $\mathbb{R}^d$ which contains incoherent $x$ with $\ell_2$ norm bounded by 1, that is, we consider,

$$\mathcal{B} = \left\{ x : \|x\|_\infty \leq \frac{\mu}{\sqrt{d}}, \|x\| \leq 1 \right\}.$$  \hspace{1cm} (3.5)

Note that the desired solution $z$ is in $\mathcal{B}$, and the regularization $R(x)$ vanishes inside $\mathcal{B}$.

The general rational is to assume $x$ satisfies the first and second order optimality condition, and then use these conditions to deduce a sequence of properties that $x$ must satisfy.

**Lemma 3.4.** Under the setting of Theorem 3.2, with high probability over the choice of $\Omega$, for any $x \in \mathcal{B}$ that satisfies second-order optimality condition (3.3) we have,

$$\|x\|^2 \geq 1/4.$$

The same is true if $x \in \mathcal{B}$ only satisfies $\tau$-relaxed second order optimality condition for $\tau \leq 0.1p$.

**Proof.** We plug in $v = z$ in the second-order optimality condition (3.3), and obtain that

$$\|P_\Omega(zz^T + xx^T)\|_F^2 \geq 2\tau^2 P_\Omega(M - xx^T)zz^T.$$ \hspace{1cm} (3.6)

Intuitively, when restricted to $\Omega$, the squared Frobenius on the LHS and the quadratic form on the RHS should both be approximately a $p$ fraction of the unrestricted case. In fact, both LHS and RHS can be written as the sum of terms of the form ($P_\Omega(vv^T)$, $P_\Omega(st^T)$), because

$$\|P_\Omega(zz^T + xx^T)\|_F^2 = 2(P_\Omega(zz^T), P_\Omega(zz^T)) + 2(P_\Omega(zz^T), P_\Omega(xx^T)) + 2(P_\Omega(xx^T), P_\Omega(zz^T)) = 2\tau^2 P_\Omega(M - xx^T)zz^T.$$ \hspace{1cm} (3.6)

Therefore we can use concentration inequalities (Theorem D.1), and simplify the equation

LHS of (3.6) = $p \|xx^T + zz^T\|_F^2 + O(\sqrt{pd}\|x\|\|z\|\|x\|\|z\|\|\|z\|^2})$

$= 2p\|x\|^2\|z\|^2 + 2p(\|x\|\|z\|^2 + O(p\epsilon), \hspace{1cm} (\text{Since } x, z \in \mathcal{B})$
where $\varepsilon = \mathcal{O}(\mu^2 \sqrt{\frac{\log d}{pd}})$. Similarly, by Theorem D.1, again, we have
\[
\text{RHS of (3.6)} = 2 \left( \langle P_{\Omega}(zz^T), P_{\Omega}(zz^T) \rangle - \langle P_{\Omega}(xx^T), P_{\Omega}(zz^T) \rangle \right)
= 2p\|z\|^4 - 2p(x,z)^2 \pm \mathcal{O}(pe)
\]
(Since $M = zz^T$)
(by Theorem D.1 and $x,z \in \mathcal{B}$)

(Note that even we use the $\tau$-relaxed second order optimality condition, the RHS only becomes $1.99p\|z\|^4 - 2p(x,z)^2 \pm \mathcal{O}(pe)$ which does not effect the later proofs.)

Therefore plugging in estimates above back into equation (3.6), we complete the proof. □

Next we use first order optimality condition to pin down another property of $x$ – it has to be close to $z$ after scaling. Note that this doesn’t mean directly that $x$ has to be close to $z$ since $x=0$ also satisfies first order optimality condition (and therefore the conclusion (3.7) below).

**Lemma 3.5.** With high probability over the randomness of $\Omega$, for any $x \in \mathcal{B}$ that satisfies first-order optimality condition (3.2), we have that $x$ also satisfies
\[
\|\langle z, x \rangle z - \|x\|^2 x\| \leq \mathcal{O}(\varepsilon).
\]
where $\varepsilon = \tilde{O}(\mu^3(pd)^{-1/2})$.

**Proof.** Note that since $x \in \mathcal{B}$, we have $R(x) = 0$. Therefore first-order optimality condition says that
\[
P_{\Omega}(M - xx^T)x = P_{\Omega}(zz^T)x - P_{\Omega}(xx^T)x = 0.
\]
Again, intuitively we hope $P_{\Omega}(zz^T) \approx pzz^T$ and $P_{\Omega}(xx^T) \approx p\|x\|^2 x$. These are made precise by the concentration inequalities Lemma D.4 and Theorem D.2 respectively.

By Theorem D.2, we have that with high probability over the choice of $\Omega$, for every $x \in \mathcal{B}$,
\[
\|P_{\Omega}(xx^T)x - pxx^T x\|_F \leq pe\|x\|^3 \leq pe
\]
where $\varepsilon = \tilde{O}(\mu^3(pd)^{-1/2})$. Similarly, by Lemma D.4, we have that for with high probability over the choice of $\Omega$,
\[
\|P_{\Omega}(zz^T) - pzz^T\| \leq \varepsilon p.
\]
for $\varepsilon = \tilde{O}(\mu^2(pd)^{-1/2})$. Therefore for every $x$,
\[
\|P_{\Omega}(zz^T)x - pzz^T x\| \leq \varepsilon p\|x\| \leq \varepsilon p.
\]
Plugging in estimates (3.10) and (3.9) into equation (3.8), we complete the proof. □

Finally we combine the two optimality conditions and show (3.7) implies $xx^T$ must be close to $zz^T$.

**Lemma 3.6.** Suppose vector $x$ satisfies that $\|x\|^2 \geq 1/4$, and that $\|\langle z, x \rangle z - \|x\|^2 x\| \leq \delta$. Then for $\delta \in (0, 0.1)$,
\[
\|xx^T - zz^T\|_F \leq \mathcal{O}(\delta).
\]

**Proof.** We write $z = ux + v$ where $u \in \mathcal{R}$ and $v$ is a vector orthogonal to $x$. Now we know $\langle z, x \rangle z = u^2\|x\|^2 x + u\|x\|^2 v$, therefore
\[
\delta \geq \|\langle z, x \rangle z - \|x\|^2 x\| = \|x\|^2 \sqrt{u^2\|v\|^2 + (1 - u^2)^2}.
\]
In particular, we know $|1 - u^2| \leq 4\delta$ and $u\|v\| \leq 4\delta$. This means $|u| \in [1 - 3\delta]$ and $\|v\| \leq 8\delta$. Now we expand $xx^T - zz^T$:
\[
xx^T - zz^T = (1 - u^2)xx^T + uvx^T + vvx^T + vv^T.
\]
It is clear that all the terms have norm bounded by $O(\delta)$, therefore $\|xx^T - zz^T\|_F \leq O(\delta)$. □
Figure 1: Partition of $\mathbb{R}^d$ into regions where our Lemmas apply. For example, Lemma 3.8 rules out the possibility that a point $x$ in the green region is local minimum. Here, the green region is the intersection of $\ell_\infty$ norm ball and $\ell_2$ norm ball. Both the white region and yellow region have non-zero gradient but for different reasons.

3.2 Extension to general $x$

We have shown when $x$ is incoherent and satisfies first and second order optimality conditions, then it must be close to $z$ or $-z$. Now we need to consider more general cases when $x$ may have some very large coordinates. Here the main intuition is that the first order optimality condition with a proper regularizer is enough to guarantee that $x$ cannot have an entry that is too much bigger than $\mu / \sqrt{d}$.

Lemma 3.7. With high probability over the choice of $\Omega$, for any $x$ that satisfies first-order optimality condition (3.2), we have

$$\|x\|_{\infty} \leq 4 \max \{\alpha, \mu \sqrt{p/\lambda} \}.$$  

(3.11)

Here we recall that $\alpha$ was chosen to be $10\mu / \sqrt{d}$ and $\lambda$ is chosen to be large so that the $\alpha$ dominates the second term $\mu \sqrt{p/\lambda}$ in the setting of Theorem 3.2.

Proof of Lemma 3.7. Suppose $i^* = \max_j |x_j|$. Without loss of generality, suppose $x_{i^*} \geq 0$. Suppose $i^*$-th row of $\Omega$ consists of entries with index $[i] \times S_{i^*}$. If $|x_{i^*}| \leq 2\alpha$, we are done. Therefore in the rest of the proof we assume $|x_{i^*}| > 2\alpha$.

Note that when $p \geq c (\log d)/d$ for sufficiently large constant $c$, with high probability over the choice of $\Omega$, we have $|S_{i^*}| \leq 2pd$. In the rest of argument we are working with such an $\Omega$ with $|S_{i^*}| \leq 2pd$.

We will compare the $i^*$-th coordinate of LHS and RHS of first-order optimality condition (3.2). For preparation, we have

$$|P_{i^*}(M)x_{i^*}| = \left| \left[ P_{[i]}(zz^T)x \right]_{i^*} \right| = \sum_{j \in S_{i^*}} |z_{i^*}z_jx_j| \leq |x_{i^*}| \sum_{j \in S_{i^*}} |z_{i^*}z_j| \leq |x_{i^*}| \cdot \mu^2/d \cdot |S_{i^*}| \leq 2|x_{i^*}|p\mu^2$$  

(3.12)

where the last step we used the fact that $|S_{i^*}| \leq 2pd$. Moreover, we have that

$$(P_{i^*}(xx^T)x)_{i^*} = \sum_{j \in S_{i^*}} x_jx_j^2 \geq 0,$$

and that

$$(\lambda \nabla R(x))_{i^*} = 4\lambda (|x_{i^*}| - \alpha)^3 \text{sign}(x_{i^*}) \geq \frac{\lambda}{2} |x_{i^*}|^3$$  

(Since $x_{i^*} \geq 2\alpha)$
Now plugging in the bounds above into the $i^*$-th coordinate of equation (3.2), we obtain
\[ 4|x_{i^*}|p\mu^2 \geq 2(P_{\Omega}(M - xx^T)x)_{i^*} \geq (\lambda \nabla R(x))_{i^*} \geq \frac{\lambda}{2}|x_{i^*}|^3, \]
which implies that $|x_{i^*}| \leq 4 \sqrt{p\mu^2/\lambda}$. □

Setting $\lambda \geq \mu^2 p/\alpha^2$ and $\alpha = 10 \mu \sqrt{1/d}$, Lemma 3.7 ensures that any $x$ that satisfies first-order optimality condition is the following ball,
\[ \mathcal{B}' = \left\{ x \in \mathbb{R}^d : \|x\|_{\infty} \leq 4\alpha \right\}. \]

Then we would like to continue to use arguments similar to Lemma 3.4 and Lemma 3.5. However, things have become more complicated as now we need to consider the contribution of the regularizer.

**Lemma 3.8** (Extension of Lemma 3.4). In the setting of Theorem 3.2 with high probability over the choice of $\Omega$, suppose $x \in \mathcal{B}'$ satisfies second-order optimality condition (3.3) or $\tau$-relaxed condition for $\tau \leq 0.1p$, we have $\|x\|^2 \geq 1/8$.

The guarantees and proofs are very similar to Lemma 3.4. The main intuition is that we can restrict our attentions to coordinates whose regularizer is equal to 0. See Section A for details.

We will now deal with first order optimality condition. We first write out the basic extension of Lemma 3.5, which follows from the same proof except we now include the regularizer term.

**Lemma 3.9** (Basic extension of Lemma 3.5). With high probability over the randomness of $\Omega$, for any $x \in \mathcal{B}'$ that satisfies first-order optimality condition (3.2), we have that $x$ also satisfies
\[ \|\langle z, x \rangle z - \|x\|^2 x - \gamma \cdot \nabla R(x)\| \leq O(\epsilon). \] (3.13)
where $\epsilon = \tilde{O}(\mu^p(pd)^{-1/2})$ and $\gamma = \lambda/(2p) \geq 0$.

Next we will show that we can remove the regularizer term, the main observation here is nonzero entries $\nabla R(x)$ all have the same sign as the corresponding entries in $x$. See Section A for details.

**Lemma 3.10.** Suppose $x \in \mathcal{B}'$ satisfies that $\|x\|^2 \geq 1/8$, under the same assumption as Lemma 3.9 we have,
\[ \|\langle x, z \rangle z - \|x\|^2 x\| \leq O(\epsilon) \] (3.14)

Finally we combine Lemma 3.7, Lemma 3.8 and Lemma 3.10 and prove Lemma 3.6. The argument are also summarized in Figure 1 where we partition $\mathbb{R}^d$ into regions where our lemmas apply.

### 4 Rank-r case

In this section we show how to extend the results in Section 3 to recover matrices of rank $r$.

Recall that in this case we assume the original matrix $M = ZZ^T$, where $Z \in \mathbb{R}^{d \times r}$. We also assume Assumption 1. The objective function is very similar to the rank 1 case
\[ f(X) = \frac{1}{2} \|P_{\Omega}(M - XX^T)\|_{F}^2 + \lambda R(X), \] (4.1)
where $R(X) = \sum_{i=1}^d r(||X||)$. and recall that $r(t) = (|t| - \alpha)^+ I_{[\alpha, \infty]}$. Here $\alpha$ and $\lambda$ are again parameters that we will determined later.

Without loss of generality, we assume that $\|Z\|^2 = r$ in this section. This implies that $\sigma_{\text{max}}(Z) \geq 1 \geq \sigma_{\text{min}}(Z)$.

Now we shall state the first and second order optimality conditions:
Proposition 4.1. If $X$ is a local optimum of objective function (4.1), its first order optimality condition is,
\[
2P_Ω(M)X = 2P_Ω(XX^T)X + \lambda \nabla R(X),
\]
and the second order optimality condition is equivalent to
\[
\forall V \in \mathbb{R}^{d \times r}, \quad \|P_Ω(VX^T + XV^T)\|_F^2 + \lambda \langle V^T, \nabla^2 R(X)V \rangle \geq 2\langle P_Ω(M - XX^T), VV^T \rangle.
\]

Note that the regularizer now is more complicated than the one dimensional case, but luckily we still have the following nice property.

Proposition 4.2. We have that $\nabla R(X) = \Gamma X$ where $\Gamma \in \mathbb{R}^{d \times d}$ is a diagonal matrix with $\Gamma_i = \frac{4d\|X_i\| - \alpha^4}{\|X_i\| - \|X_i\|}_2 \|x_i\|$. As a direct consequence, $\langle (\nabla R(X))_i, X_i \rangle = 0$ for every $i \in [d]$.

Now we are ready to state the precise version of Theorem 1.1:

Theorem 4.3. Suppose $p \geq C \max\{\mu / \kappa^6 r^4, \mu / \kappa^4 r^4\}d^{-1} \log^{1.5} d$ where $C$ is a large enough constant. Let $\alpha = 4\mu r / \sqrt{d}, \lambda \geq \mu^2 r p / \alpha^2$. Then with high probability over the randomness of $\Omega$, any local minimum $X$ of $f(\cdot)$ satisfies that $f(X) = 0$, and in particular, $ZZ^T = XX^T$.

The proof of this Theorem follows from a similar path as Theorem 3.2. We first notice that because of the regularizer, any matrix $X$ that satisfies first order optimality condition must be somewhat incoherent (this is analogous to Lemma 3.7):

Lemma 4.4. Suppose $|S_1| \leq 2pd$. Then for any $X$ satisfies 1st order optimality condition (4.2), we have
\[
\|X\|_{2 \rightarrow \infty} = \max_i \|X_i\| \leq 4 \max\{\alpha, \mu \sqrt{rp}/\lambda\}
\]

Proof. Assume $i^* = \arg\max_j \|X_i\|$. Suppose the $i$th row of $\Omega$ consists of entries with index $[i] \times S_1$. If $\|X_i\| \leq 2\alpha$, then we are done. Therefore in the rest of the proof we assume $\|X_i\| \geq 2\alpha$.

We will compare the $i$-th row of LHS and RHS of (4.2). For preparation, we have
\[
(P_Ω(M)x)_i = (P_Ω(ZZ^T))_i = (P_Ω(ZZ^T))_i X
\]
Then we have that
\[
\left\|\left(\frac{(P_Ω(ZZ^T))_i}{X_i}\right)\right\|_1 = \sum_{j \in S_1} \|Z_{i, j}\| \leq \sum_{j \in S_1} \|Z_{i, j}\| \leq \sum_{j \in S_1} \mu^2 r/|S_1| \leq 2\mu^2 r p (by \text{incoherence of } Z)
\]
\[
= 2\mu^2 r p (by \|S_1\| \leq 2pd)
\]
Therefore we can bound the $\ell_2$ norm of LHS of 1st order optimality condition (4.2) by
\[
\left\|\left(\frac{(P_Ω(ZZ^T)X)_i}{X_i}\right)\right\|_2 \leq \left\|\left(\frac{(P_Ω(ZZ^T))_i}{X_i}\right)\right\|_1 \|X_i\|_1 \leq 2\mu^2 r p \|X_i\|_2 \rightarrow \infty (by \|X_i\|_{2 \rightarrow \infty} = \|X_i\|_1 \rightarrow 2)
\]
\[
= 2\mu^2 r p \|X_i\| (4.6)
\]
Next we lowerbound the norm of the RHS of equation (4.2). We have that
\[
(P_Ω(XX^T)X)_i = \sum_{j \in S_1} \langle X_{i, j}, X_j \rangle X_j = X_i \sum_{j \in S_1} X_j^T X_j,
\]
which implies that
\[
\langle (P_Ω(XX^T))_i, X_i \rangle = X_i \left\{ \sum_{j \in S_1} X_j^T X_j \right\} X_i^T \geq 0. \quad (4.7)
\]
Using Proposition 4.2, we obtain that
\[ \langle (P_\Omega(XX^T)X)_r, (\nabla R(X))_r \rangle = \Gamma_\mu X_r \left( \sum_{j \in X_r} X_j^T X_j \right) X_r^\top \geq 0. \] (4.8)

It follows that
\[ \left\| (P_\Omega(XX^T)X)_r + (\lambda \nabla R(X))_r \right\| \geq \left\| (\lambda \nabla R(X))_r \right\| \] (by equation (4.8))
\[ = \frac{4 \lambda (\| X_r \| - \delta)^3}{\| X_r \|} \cdot \| X_r \| \] (by Proposition 4.2)
\[ \geq \frac{\lambda}{2} \| X_r \|^3 \] (by the assumption \| X_r \| \geq 2 \alpha)

Therefore plugging in equation above and equation (4.6) into 1st order optimality condition (4.2). We obtain that \| X_r \| \leq \sqrt{8 \mu r p / \lambda} which completes the proof. \( \square \)

Next, we prove a property implied by first order optimality condition, which is similar to Lemma 3.9

**Lemma 4.5.** In the setting of Theorem 4.3 with high probability over the choice of \( \Omega \), for any \( X \) that satisfies 1st order optimality condition (4.2), we have
\[ \| X \|^2_F \leq 2 r \sigma_{\max}(Z)^2. \] (4.9)
Moreover, we have
\[ \sigma_{\max}(X) \leq 2 \alpha \sigma_{\max}(Z)^{-1/6}. \] (4.10)
and
\[ \| ZZ^T X - XX^T Y - \gamma \nabla R(X) \|_F \leq O(\delta) \] (4.11)
where \( \delta = O(\mu^3 \kappa^3 \gamma^2 \log^{0.75}(d) \sigma_{\max}(Z)^{-3}(dp)^{-1/2}) \) and \( \gamma = \lambda / (2p) > 0 \).

**Proof.** If \( \| X \|_F \leq \sqrt{r \sigma_{\max}(Z)^2} \) we are done. When \( \| X \|_F \geq \sqrt{r \sigma_{\max}(Z)^2} \), by Lemma 4.4, we have that max \( \| X \| \leq 4 \alpha = 4 \mu \kappa r / \sqrt{d} \), and therefore max \( \| X \| \leq \nu \| X \|_F \) with \( \nu = O(\mu \kappa r^2 / \sigma_{\max}(Z)) \). Then by Theorem D.2, we have that
\[ \left\| P_\Omega(ZZ^T X - p ZZ^T X) \right\|_F \leq p \delta, \] (4.12)
and
\[ \left\| P_\Omega(XX^T X - XX^T X) \right\|_F \leq p \delta, \] (4.13)
where \( \delta = O(\mu^3 \kappa^3 \gamma^2 \log^{0.75}(d) \sigma_{\max}(Z)^{-3}(dp)^{-1/2}) \). These two imply equation (4.11). Moreover, we have
\[ p \| ZZ^T X \|_F = \left\| P_\Omega(ZZ^T X) \right\|_F + p \delta = \left\| P_\Omega(XX^T X + \lambda R(X)) \right\|_F + p \delta \] (by equation (4.2))
\[ \geq \left\| P_\Omega(XX^T X) \right\|_F + p \delta \] (by equation (4.8))
\[ \geq p \| XX^T X \|_F + 2p \delta \] (4.14)

Suppose \( X \) has singular value \( \sigma_1 \geq \ldots \geq \sigma_r \). Then we have \( \| ZZ^T X \|_F^2 \leq \| ZZ^T \|^2 \| X \|^2_F \leq \sigma_{\max}(Z)^4 \| X \|^2_F = \sigma_{\max}(Z)^4 (\sigma_1^2 + \cdots + \sigma_r^2) \). On the other hand, \( \| XX^T X \|_F^2 = \sigma_1^6 + \cdots + \sigma_r^6 \). Therefore, equation (4.14) implies that
\[ (1 + O(\delta)) \sigma_{\max}(Z)^4 \sum_{i=1}^r \sigma_i^2 \geq \sum_{i=1}^r \sigma_i^6 \] (4.15)

Then we have (by Proposition E.1) we complete the proof. \( \square \)
Now we look at the second order optimality condition, this condition implies the smallest singular value of $X$ is large (similar to Lemma 3.8). Note that this lemma is also true even if $x$ only satisfies relaxed second order optimality condition with $\tau = 0.01 \sigma_{\min}(Z)$.

**Lemma 4.6.** In the setting of Theorem 4.3, with high probability over the choice of $\Omega$, suppose $X$ satisfies equation (4.4), the second order optimality condition (4.3). Then,

$$\sigma_{\min}(X) \geq \frac{1}{4} \sigma_{\min}(Z) \quad (4.16)$$

**Proof.** Let $J = \{i : \|X_i\| \leq \alpha\}$. Let $v \in \mathbb{R}^d$ such that $\|Xv\| = \sigma_{\min}(X)$. Let $Z_J$ be the matrix that has the same $i$-th row as $Z$ for every $i \in J$ and 0 elsewhere. Since $Z_J$ has column rank at most $r$, by variational characterization of singular values, we have that for there exists unit vector $z \in \text{col-span}(Z_J)$ such that $\|X^T z\| \leq \sigma_{\min}(X)$.

We claim that $\sigma_{\min}(Z_J) \geq \frac{1}{4} \sigma_{\min}(Z)$. Let $L = [d] - J$. Since for any $i \in L$ it holds that $\|X_i\| \geq \alpha$, we have $|\langle L \rangle^2 | \leq 2r \sigma_{\max}(Z)^2$ (by equation (4.9)), and it follows that $|\langle L \rangle | \leq 2r \sigma_{\max}(Z)^2 / \alpha^2$. Therefore,

$$\sigma_{\min}(Z_J) \geq \sigma_{\min}(Z) - \sigma_{\max}(Z_J) \geq \sigma_{\min}(Z) - \|Z_J\|_F$$

$$\geq \sigma_{\min}(Z) - \sqrt{|\langle L \rangle |^2 \delta} \geq \sigma_{\min}(Z) - \sqrt{\varepsilon^2 \sigma_{\max}(Z)^2 / \alpha^2} \geq \frac{1}{5} \sigma_{\min}(Z). \quad (\text{by } \alpha \geq \frac{\varepsilon}{\sqrt{d}})$$

Since $z \in \text{col-span}(Z_J)$ is a unit vector, we have that $z_J = Z_J \beta$ where $\|\beta\| \leq \frac{1}{\sqrt{\sigma_{\min}(Z_J)}} \leq O(1/\sigma_{\min}(Z))$. Therefore this in turn implies that $\|z_J\|_\infty \leq \|Z_J\|_{2 \rightarrow \infty} \|\beta\| \leq O(\mu \sqrt{r/d}) / \sigma_{\min}(Z) \leq O(\mu \sqrt{r/d})$.

We will plug in $V = z_J v^T$ in the second order optimality condition (4.3). Note that since $z_J \in \text{col-span}(Z_J)$, it is supported on subset $J$, and therefore $\nabla^2 R(X) V = 0$. Therefore the term about regularization in (4.3) will vanish. For simplicity, let $y = X^T z_J$, $w = Xv$. We obtain that taking $V = z_J v^T$ in equation (4.3) will result in

$$\|P_\Omega(w z_J^T + z_J w^T)\|_F^2 \geq 2\langle P_\Omega(Z Z^T - XX^T), z_J z_J^T \rangle \quad (4.17)$$

Note that we have that $\|w\|_\infty \leq \|X\|_{2 \rightarrow \infty} |v| \leq \mu \sqrt{r/d}$. Recalling that $\|z_J\|_\infty \leq O(\mu \sqrt{r/d})$, by Theorem D.1, we have that

$$p \|w z_J^T + z_J w^T\|_F^2 \geq 2p(\|Z Z^T - XX^T, z_J z_J^T\| - \delta p) \quad (4.18)$$

where $\delta = O(\mu^2 \kappa r^2 (pd)^{-1/2})$. Then simple algebraic manipulation gives that

$$\langle w, z_J \rangle^2 + \|w\|^2 \|z_J\|^2 + \|X^T z_J\|^2 \geq \|Z^T z_J\|^2 - \delta / 2 \quad (4.19)$$

Note that $\langle w, z_J \rangle = \langle v, X^T z_J \rangle = \langle v, y \rangle$. Recall that $|\langle z_J \rangle| = 1$ and $z \in \text{col-span}(Z)$, and therefore $\|Z^T z_J\| = \|Z_J^T z_J\| \geq \sigma_{\min}(Z_J)$, Moreover, recall that $\|w\| = \|X^T z_J\| \leq \sigma_{\min}(X)$. Using these with equation (4.19) we obtain that

$$\langle w, z_J \rangle^2 + \|w\|^2 \|z_J\|^2 + \|X^T z_J\|^2 \leq \langle y, v \rangle^2 + \|y\|^2 + \|v\|^2 \leq 2\|v\|^2 + \sigma_{\min}(X) \quad (\text{by Cauchy-Schwarz and } \|w\| = \sigma_{\min}(X).$$

$$\leq 3\sigma_{\min}(X) \quad (\text{by } \|v\| \leq \sigma_{\min}(X).)$$

Therefore together with equation (4.19) and $\|Z^T z_J\| \geq \sigma_{\min}(Z_J)$ we obtain that

$$\sigma_{\min}(X) \geq (1/2 - \Omega(\delta))\sigma_{\min}(Z_J) \quad (4.20)$$

Therefore combining equation (4.20) and the lower bound on $\sigma_{\min}(Z_J)$ we complete the proof. \qed

Similar as before, we show it is possible to remove the regularizer term here, again the intuition is that the regularizer is always in the same direction as $X$. 

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Lemma 4.7. Suppose $X$ satisfies equation (4.4) and (4.16) and (4.10), then for any $\gamma \geq 0$,\

$$\| Z Z^T X - X X^T X \|_F^2 \leq \| Z Z^T X - X X^T X - \gamma \nabla R(X) \|_F^2 \quad (4.21)$$

Proof. Let $L = \{ i : \| X_i \|_2 \geq \alpha \}$. For $i \notin L$, we have that $(\nabla R(X))_i = 0$. Therefore it suffices to prove that for every $i \in L$,

$$\| Z_i Z_i^T X - X_i X_i^T X \|_F^2 \leq \| Z_i Z_i^T X - X_i X_i^T X - (\gamma \nabla R(X))_i \|_F^2$$

It suffices to prove that

$$\langle (\nabla R(X))_i, X_i X_i^T X - Z_i Z_i^T Z \rangle \geq 0 \quad (4.22)$$

By proposition 4.2 we have $\nabla R(X)_i = \Gamma X_i$ for $\Gamma \geq 0$. Then

$$\langle (\nabla R(X))_i, X_i X_i^T X \rangle = \langle \Gamma X_i, X_i X_i^T X \rangle \geq \Gamma \| X_i \|_2^2 \sigma_{\min}(X_i X_i^T) \geq \frac{1}{4} \Gamma \| X_i \|_2^2 \sigma_{\min}(Z) \quad \text{(by equation 4.16)}$$

On the other hand, we have

$$\langle (\nabla R(X))_i, Z_i Z_i^T Z \rangle = \langle \Gamma X_i, Z_i Z_i^T X \rangle$$

$$\leq \Gamma \| X_i \|_2 \| Z_i \|_2 \| Z_i \|_2 \sigma_{\max}(Z_i Z_i^T) \leq \Gamma \| X_i \|_2 \| Z_i \|_2 \sigma_{\max}(Z_i Z_i^T) \sigma_{\max}(X_i X_i^T) \leq \Gamma \| X_i \|_2 \| Z_i \|_2 \sigma_{\max}(Z_i Z_i^T)^2 r_{1/6} \quad \text{(by equation 4.10)}$$

$$\leq \frac{1}{10} \Gamma \| X_i \|_2^2 \sigma_{\max}(Z_i Z_i^T)^2 r^{-1/3} \quad \text{(by } \| X_i \|_2 \geq \alpha \geq 10 \sqrt{r} \| X_i \|)$$

Therefore combining two equations above we obtain equation (4.22) which completes the proof. \qed

Finally we show the form in Equation (4.21) implies $Z Z^T$ is close to $X X^T$ (this is similar to Lemma 3.6).

Lemma 4.8. Suppose $X$ and $Z$ satisfies that $\sigma_{\min}(X) \geq 1/4 \cdot \sigma_{\min}(Z)$ and that

$$\| Z Z^T X - X X^T X \|_F^2 \leq \delta^2$$

(4.23)

where $\delta \leq \sigma_{\min}(Z)/C$ for a large enough constant $C$, then

$$\| X X^T - Z Z^T \|_F^2 \leq O(\delta^2 \sigma_{\min}(Z)).$$

(4.24)

Proof. The proof is similar to the one-dimensional case, we will separate $Z$ into the directions that are in column span of $X$ and its orthogonal subspace. We will then show the projection of $Z$ in the column span is close to $X$, and the projection on the orthogonal subspace must be small.

Let $Z = U + V$ where $U = \text{Proj}_{\text{span}(X)} Z$ is the projection of $Z$ to the column span of $X$, and $V$ is the projection to the orthogonal subspace. Then since $V X = 0$ we know

$$Z Z^T X = (U + V)(U + V)^T X = U U^T X + V V^T X.$$ 

Here columns of the first term $U U^T X$ are in the column span of $X$, and the columns second term $V V^T X$ are in the orthogonal subspace. Therefore,

$$\| Z Z^T X - X X^T X \|_F^2 = \| U U^T X - X X^T X \|_F^2 + \| V V^T X \|_F^2 \leq \delta^2.$$ 

In particular, both terms should be bounded by $\delta^2$. Therefore $\| U U^T X - X X^T X \|_F^2 \leq \delta^2 \sigma_{\min}(X)^2 \leq 16 \delta^2 \sigma_{\min}(Z)^2$.

Also, we know $\sigma_{\min}(U U^T X) \geq \sigma_{\min}(X X^T X) - \delta \geq \sigma_{\min}(Z)^2/128$ if $\delta \leq \sigma_{\min}(Z)^2/128$. Therefore $\sigma_{\min}(U U^T X)$ is at least $\sigma_{\min}(Z)^2/\|Z\|_2^2$. Now $\| V \|_F^2 \leq \delta^2 \sigma_{\min}(U U^T X)^2 \leq O(\delta^2 \|Z\|_2^2)$. Finally, we can bound $\| U V \|_F$ by $\| U \|_F \| V \|_F \leq \|Z\|_2 \|V\|_F$ (last inequality is because $U$ is a projection of $Z$), which at least $\Omega(\|v\|_2^2)$ when $\delta \leq \sigma_{\min}(Z)^3/128$, therefore

$$\| Z Z^T - X X^T \|_F \leq \| U U^T - X X^T \|_F + \| V V^T \|_F \leq O(\delta \|Z\|_2^2 \sigma_{\min}(Z)^3).$$

\qed
Last thing we need to prove the main theorem is a result from Sun and Luo\cite{SL15}, which shows whenever $XX^T$ is close to $ZZ^T$, the function is essentially strongly convex, and the only points that have 0 gradient are points where $XX^T = ZZ^T$, this is explained in Lemma C.1. Now we are ready to prove Theorem 4.3.

Proof of Theorem 4.3. Suppose $X$ satisfies 1st and 2nd order optimality condition. Then by Lemma 4.5 and Lemma 4.4 we have that $X$ satisfies equation (4.4), (4.9), (4.10) and (4.11). Then by Lemma 4.6 we obtain that $\sigma_{\min}(X) \geq 1/6 \cdot \sigma_{\min}(Z)$. Now by Lemma 4.7 and equation (4.11), we have that $\|ZZ^T X - XX^T X\|_F \leq \delta$ for $\delta \leq c \sigma_{\min}(Z)^3 / \kappa^2$ for sufficiently small constant $c$. Then by Lemma 4.8 we obtain that $\|ZZ^T X - XX^T X\|_F \leq c \sigma_{\min}(Z)^2$ for sufficiently small constant $c$. By Lemma C.1 in this region the only points that satisfy the first order optimality condition must satisfy $XX^T = ZZ^T$. \hfill \Box

Handling Noise. To handle noise, notice that we can only hope to get an approximate solution in presence of noise, and to get that our Lemmas only depend on concentration bounds which still apply in the noisy setting. See Section B for details.

5 Conclusions

Although the matrix completion objective is non-convex, we showed the objective function has very nice properties that ensures the local minima are also global. This property gives guarantees for many basic optimization algorithms. An important open problem is the robustness of this property under different model assumptions: Can we extend the result to handle asymmetric matrix completion? Is it possible to add weights to different entries (similar to the settings studied in \cite{LLKR15})? Can we replace the objective function with a different distance measure rather than Frobenius norm (which is related to works on 1-bit matrix sensing \cite{DPvdBW14})? We hope this framework of analyzing the geometry of objective function can be applied to other problems.

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A Omitted Proofs in Section 3

We first prove the equivalent form of the first and second order optimality conditions:

**Lemma A.1** (Proposition 3.1 restated). The first order optimality condition of objective (3.1) is

\[ 2P_\Omega(M - xx^\top)x = \lambda \nabla R(x), \]  

(A.1)

and the second order optimality condition requires:

\[ \forall v \in \mathbb{R}^d, \|P_\Omega(vx^\top + xv^\top)\|_F^2 + \lambda v^\top \nabla^2 R(x)v \geq 2v^\top P_\Omega(M - xx^\top)v. \]  

(A.2)

Moreover, The \( \tau \)-relaxed second order optimality condition requires

\[ \forall v \in \mathbb{R}^d, \|P_\Omega(vx^\top + xv^\top)\|_F^2 + \lambda v^\top \nabla^2 R(x)v \geq 2v^\top P_\Omega(M - xx^\top)v - \tau \|v\|^2. \]  

(A.3)

**Proof.** We take the Taylor’s expansion around point \( x \). Let \( \delta \) be an infinitesimal vector, we have

\[
\begin{align*}
f(x + \delta) &= \frac{1}{2} \|P_\Omega(M - (x + \delta)(x + \delta)^\top)\|_F^2 + \lambda R(x + \delta) + o(\|\delta\|^2) \\
&= \frac{1}{2} \|P_\Omega(M - xx^\top - (x\delta^\top + \delta x^\top) - \delta\delta^\top)\|_F^2 + \lambda \left( R(x) + \langle \nabla R(x), \delta \rangle + \frac{1}{2} \delta^\top \nabla^2 R(x)\delta \right) + o(\|\delta\|^2) \\
&= \frac{1}{2} \|M - xx^\top\|_F^2 + \lambda R(x) \\
&\quad - \langle P_\Omega(M - xx^\top), x\delta^\top + \delta x^\top \rangle + \langle \nabla R(x), \delta \rangle + o(\|\delta\|^2) \\
&\quad - \langle P_\Omega(M - xx^\top), \delta\delta^\top \rangle + \frac{1}{2} \|x\delta^\top + \delta x^\top\|_F^2 + \frac{1}{2} \lambda \delta^\top \nabla^2 R(x)\delta + o(\|\delta\|^2).
\end{align*}
\]

By symmetry \( \langle P_\Omega(M - xx^\top), x\delta^\top \rangle = \langle P_\Omega(M - xx^\top), \delta x^\top \rangle = \langle P_\Omega(M - xx^\top)x, \delta \rangle \), so the first order optimality condition is \( \forall \delta, \langle -2P_\Omega(M - xx^\top)x + \lambda \nabla R(x), \delta \rangle = 0 \), which is equivalent to that \( 2P_\Omega(M - xx^\top)x = \lambda \nabla R(x) \).

The second order optimality condition says \( -\langle P_\Omega(M - xx^\top), \delta\delta^\top \rangle + \frac{1}{2} \|x\delta^\top + \delta x^\top\|_F^2 + \frac{1}{2} \lambda \delta^\top \nabla^2 R(x)\delta \geq 0 \) for every \( \delta \), which is exactly equivalent to Equation (3.3).

Next we show the full proof for the second order optimality condition:

**Lemma A.2** (Lemma 3.8 restated). In the setting of Theorem 3.2 with high probability over the choice of \( \Omega \), suppose \( x \in \mathcal{B}' \) satisfies second-order optimality condition (3.3) or \( \tau \)-relaxed condition for \( \tau \ll 0.1 \), p, we have \( \|x\|^2 \gg 1/8 \).

**Proof.** If \( \|x\| \geq 1 \), then we are done. Therefore in the rest of the proof we assume \( \|x\| \ll 1 \). The proof is very similar to Lemma 3.4. We plug in \( v = z_J \) instead into equation (3.3), where \( J = \{i : |x_i| \leq \alpha \} \). Note that \( \nabla R(z_J) \) vanishes. We plug in \( v = z_J \) in the equation (3.3) and obtain that \( x \) satisfies that

\[ \|P_\Omega(z_J x^\top + z_J z_J^\top)\|_F^2 \geq 2z_J^\top P_\Omega(M - xx^\top)z_J. \]  

(A.4)

Note that we assume \( |x|_\infty \ll 2\alpha \), and in the beginning of the proof we assume \( \log |x| \ll 1 \). Moreover, we have \( \|z_J\| \ll \frac{\alpha}{\sqrt{\tauN}} \) an, \( \|z_J\| < 1 \). Similarly to the derivation in the proof of Lemma 3.4, we apply Theorem D.1 (twice) and obtain that with high probability over the choice of \( \Omega \), for every \( x \), for \( \varepsilon = O(\mu^2(pd)^{-\tau/2}) \),

\[
\begin{align*}
\text{LHS of (A.4)} &= p \|z_J x^\top + x_J z_J^\top\|_F^2 + O(p\varepsilon) = 2p\|x\|^2|z_J|^2 + 2p(x, z_J)^2 + O(p\varepsilon). \\
\text{RHS of (A.4)} &= 2 \left( \langle P_\Omega(zz^\top), P_\Omega(z_J z_J^\top) \rangle - \langle P_\Omega(xz^\top), P_\Omega(z_J z_J^\top) \rangle \right) \\
&= 2\|z_J\|^2 - 2(x, z_J)^2 + O(p\varepsilon). \quad \text{(by Theorem D.1)}
\end{align*}
\]
(Again notice that using $\tau$-relaxed second order optimality condition does not effect the RHS by too much, so it does not change later steps.) Therefore plugging the estimates above back into equation (A.4), we have that
\[ p\|x\|^2\|z_J\|^2 + 2p(x, z_J)^2 \geq p\|z_J\|^4 + O(\varepsilon). \]

Using Cauchy-Schwarz, we have $\|x\|^2\|z_J\|^2 \geq (x, z_J)^2$, and therefore we obtain that $\|z_J\|^2\|x\|^2 \geq \frac{1}{2}\|z_J\|^4 - O(\varepsilon)$.

Finally, we claim that $\|z_J\|^2 \geq 1/2$, which completes the proof since $\|x\|^2 \geq \frac{1}{2}\|z_J\|^2 - O(\varepsilon) \geq 1/8$.

**Claim A.3.** Suppose $\alpha > \frac{4\mu}{\nu d}$ and $x$ satisfies $\|x\|_\infty \leq 4\alpha$ and $\|x\| \leq 2$. Let $J = \{i : |x_i| < \alpha\}$. Then we have that $\|z_J\| \geq 1/2$.

The claim can be simply proved as follows: Since $\|x\|^2 \leq 2$ we have that $|J| < 2/\alpha^2$ and therefore $\|z_J\|^2 \geq 2\mu^2/(d\alpha^2)$. This further implies that $\|z_J\|^2 = \|z\|^2 - \|z_L\|^2 \geq (1 - 2\mu^2/(d\alpha^2)) > \frac{1}{2}$ because $\alpha > \frac{2\mu}{\sqrt{d}}$.

**Lemma A.4 (Lemma 3.10 restated).** Suppose $x \in B'$ satisfies that $\|x\|^2 \geq 1/8$, under the same assumption as Lemma 3.9 we have,
\[ \left\|(x, z)z - \|x\|^2x \right\| \leq O(\varepsilon) \quad \text{(A.5)} \]

**Proof.** We consider the subset of coordinates $J = \{i : |x_i| > \alpha\}$ where the regularization $g(x)$ doesn’t have any effect. Let $L = \{i : |x_i| \leq \alpha\}$ be the complement of $J$. By Claim A.3, we have that $\|z_L\| \geq 1/2$.

For coordinates that are in $J$, by equation (3.13), we have that
\[ \left\|(x, z)z - \|x\|^2x \right\| = \left\||(x, z)z - \left((\|x\|^2x_J + \gamma\nabla R(x))\right)_J\right\| \leq O(\varepsilon). \quad \text{(A.6)} \]

Note by triangle inequality, (A.6) also implies $\langle x, z \rangle \|x\|^3 - O(\varepsilon) \geq 1/20$. Therefore we can divide both side by $\langle x, z \rangle$, and let $\beta = \|x\|^2/\langle x, z \rangle$.

\[ \|z_J - \beta x_J\| \leq O(\varepsilon) \quad \text{(A.7)} \]

We now claim $\beta$ is large:

**Claim A.5.** $\beta \geq 1 - O(\varepsilon)$

**Proof.** We first claim $\|x\| \leq 1 + O(\varepsilon)$. This is because $\langle \nabla R(x), x \rangle > 0$, so $\|x\|^2 + \gamma\nabla R(x)\| \geq \|x\|^2x \| = \|x\|^3$, and $\|(x, z)z\| \leq \|x\|$. When $\|x\| \geq 1 + C\varepsilon$ for large $C$ we have $\left\||(z, x)z - \|x\|^2x - \gamma \cdot \nabla R(x)\right\| \geq C'\varepsilon$ for large $C'$ and that contradicts with (3.13).

This implies the norm of $x_J$ is bounded by $(1 + O(\varepsilon))\|z_J\|$ because for any coordinate $i$ in $L$, we know $|x_i| > \alpha > |z_i|$. Now we use (A.7) and triangle inequality:
\[ \|\beta x_J\| \geq \|z_J\| - O(\varepsilon) \geq (1 - O(\varepsilon))\|z_J\| \geq (1 - O(\varepsilon))\|x_J\| \]

Using the Claim, we can now consider the case when $i \in L$. Without loss of generality let’s assume $x_i > \alpha$. Since we have shown that $\beta x_i > (1 - O(\varepsilon))\alpha$, we have that $\beta x_i > (1 - O(\varepsilon))\alpha \geq |z_i|$. Moreover we have $\lambda\nabla R(x)_i > 0$ when $x_i > \alpha$. This implies $\|\beta x_i + \gamma\nabla R(x)_i\| \geq |z_i|$. Therefore,
\[ \|z_L - \beta x_L\| \leq \|z_L - (\beta x + \gamma\nabla R(x))_L\| \]

Combining this with (A.6), we have the result.
B Handling Noise

Suppose instead of observing the matrix $ZZ^T$, we actually observe a noisy version $M = ZZ^T + N$, where $N$ is a Gaussian matrix with independent $N(0, \sigma^2)$ entries. In this case we should not hope to exactly recover $ZZ^T$ (as two close $Z$’s may generate the same observation). In this Section we show even with fairly large noise our arguments can still hold.

**Theorem B.1.** Let $\hat{\mu} = \max[\mu, \sqrt{4rd \log d / r}]$. Suppose $p \geq C^2 \mu^2 k^{12} r d^{-1} e^{-2} \log^2 d$ where $C$ is a large enough constant. Let $\alpha = 2\hat{\mu} kr / \sqrt{d} \lambda \geq \mu^2 r / \alpha^2$. Then with high probability over the randomness of $\Omega$, any local minimum $X$ of $f(\cdot)$ satisfies

$$\|XX^T - ZZ^T\|_F \leq \epsilon.$$

In fact, a noise level $\sigma \sqrt{\log d} \leq \mu^2 r / d$ (when the noise is almost as large as the maximum possible entry) does not change the conclusions of Lemmas in this Section.

**Proof.** There are only three places in the proof where the noise will make a difference. These are: 1. The infinity norm bound of $M$, used in Lemma 3.2. The LHS of first order optimality condition (Equation 4.2). 3. The RHS of second order optimality condition (Equation 4.3).

What we require in these three steps are: 1. $|M|_{\infty}$ should be smaller than $\mu^2 r / d$. 2. $\langle P_{\Omega}(N), W \rangle$ should be smaller than $\langle P_{\Omega}(N), W \rangle \leq O(\sigma |Z|_{\infty} dr \log d + \sqrt{pd^2 r \log d |W|_{\infty} \| W \|_F \log d}$. 3. $\|P_{\Omega}(N)\| \leq \epsilon \|ZZ^T\|_F$. When we define the $\hat{\mu} = \max[\mu, \sqrt{4rd \log d / r}]$, all of these are satisfied (by Lemma D.5 and D.6).

Now we can follow the proof and see $\delta \leq c \sigma_{\min}(Z) / \epsilon^2$ for small enough constant $c$, and By Lemma 4.8 we know $\|XX^T - ZZ^T\|_F \leq \epsilon$. \qed

C Finding the Exact Factorization

In Section B we showed that any point that satisfies the first and second order necessary condition must satisfy $\|XX^T - ZZ^T\|_F \leq c$ for a small enough constant $c$. In this section we will show that in fact $XX^T$ must be exactly equal to $ZZ^T$. The proof technique here is mostly based on the work of Sun and Luo [SL15]. However we have to modify their proof because we use slightly different regularizers, and we work in the symmetric case. The main Lemma in [SL15] can be rephrased as follows in our setting:

**Lemma C.1** (Analog to Lemma 3.1 in [SL15]). Assuming $p$ is at least $C \mu^4 k^4 d^{-1} \log d$ for large enough $d$, for any $X$ such that $\|XX^T - ZZ^T\|_F \leq \epsilon \leq \sigma_{\min}(Z)^2 / 100$, $\|X\|_{2\rightarrow \infty}$ is bounded as in Lemma 4.4 there must be a matrix $U$ such that $UU^T = ZZ^T$ and

$$\langle \nabla f(X), X - U \rangle \geq \frac{p}{4} \|M - XX^T\|_F^2.$$

As a corollary, the points in this set that satisfy first order optimality condition must have $XX^T = ZZ^T$.

In order to prove this main lemma, we separate $f(X) = g(X) + R(X)$ where $g(X) = \frac{1}{2} \|P_{\Omega}(M - XX^T)\|_F^2$, and $R(X)$ is the regularizer. We will first show that the regularizer always has a positive correlation with $X - U$.

**Claim C.2.** For any $U$ such that $UU^T = ZZ^T$,

$$\langle \nabla R(X), X - U \rangle \geq 0.$$

**Proof.** Since the regularizer is applied to individual rows, we will focus on a row $X_i$. The key observation here is that for each row $X_i$, $\nabla R(X_i)$ is 0 when $\|X_i\| \leq 2\mu / \sqrt{d} \leq \alpha$. When $\|X_i\|$ is larger $\nabla R(X_i)$ is always in the same direction as $X_i$. Therefore we know either $\nabla R(X_i) = 0$, or $\nabla R(X_i) = \lambda X_i$ for some $\lambda > 0$ and $\|X_i\| > 2\mu / \sqrt{d} \geq 2\|Z_i\| = 2\|U_i\|$. The innerproduct $\langle \nabla R(X_i), X_i - U_i \rangle$ is 0 in the former case, and is at least $\lambda (\|X_i\|^2 - \|X_i\|\|U_i\|)$ in the latter case. \qed

Note that this Claim is much more complicated in [SL15] (Claim 4.1) in order to deal with the asymmetric case. Next we will prove the gradient of $g(X)$ has a large correlation with $X - U$.
Claim C.3. There exists a matrix \( U \) where \( UU^T = ZZ^T \), such that

\[
\langle \nabla R(X), X - U \rangle \geq \frac{P}{4} \| M - XX^T \|_F^2.
\]

The proof of this Claim follows from the same strategy as in [SL15]. We will first show that there exists a matrix \( U \) such that \( \| U - X \|_F \) is small, and then use concentration bounds to show the lowerbound on innerproduct. The concentration bounds are exactly the same as [SL15] (see the proof of Proposition 4.3), the only thing we need to do here is an equivalent version of Proposition 4.1 and 4.2:

Claim C.4. Suppose \( \| XX^T - M \|_F = \epsilon \leq \sigma_{\min}(Z)^2 / 100 \), there exists a matrix \( U \) such that \( UU^T = M \) and \( \| X - U \|_F \leq 5\epsilon \sqrt{r / \sigma_{\min}(Z)^2} \).

Note that in this claim we do not need to prove anything about the row norm/Frobenius norm of \( U \), because whenever \( UU^T = M \), \( U \) must have the same row norm and Frobenius norm as the original low rank component \( Z \). That is why the proof here is again much simpler than [SL15].

Proof. Without loss of generality we assume \( M \) is a diagonal matrix with first \( r \) diagonal terms being \( \sigma_1(Z)^2, \sigma_2(Z)^2, \ldots, \sigma_r(Z)^2 \) (this can be done by a change of basis). We use \( M' \) to denote the first \( r \times r \) submatrix of \( M \).

Now we write \( X \) in a block form \( X^T = (V^T W^T) \), where \( V \) is an \( r \times r \) matrix and \( W \) is an \((d - r) \times r \) matrix. Clearly, we want to truncate the \( W \) part (because the corresponding part in \( M \) is empty), and we need to “fix” the \( V \) part so that the first \( r \times r \) submatrix in \( XX^T \) is exactly equal to \( M \).

We will construct the first \( r \times r \) block of \( U \) as \( U = VQ = V(V^T M^{-1} V)^{-1/2} \) (where \( Q := (V^T M^{-1} V)^{-1/2} \)), all the other entries of \( U \) are 0. The difference in \( U \) and \( X \) is equal to \( \| U - X \|_F \leq \| P - V \|_F + \| W \|_F \).

Since \( \| XX^T - M \|_F = \epsilon \), we know \( \| M' - VV^T \|_F^2 + 2\| VV^T\|_F^2 \leq \epsilon^2 \). In particular both terms are smaller than \( \epsilon^2 \).

First, we bound \( \| W \|_F \). Note that since \( \| M' - VV^T \|_F \leq \epsilon \leq \sigma_{\min}(Z)^2 / 100 \), we know \( \sigma_{\min}(V)^2 \geq 0.9\sigma_{\min}(Z)^2 \). Therefore \( \sigma_{\min}(V) \geq 0.9\sigma_{\min}(Z) \). Now we know \( \| W \|_F \leq \| WW^T \|_F / \sigma_{\min}(V) \leq 2\epsilon / \sigma_{\min}(Z) \).

Next we bound \( \| P - V \|_F^2 \). Since \( \| M' - VV^T \|_F \leq \epsilon \leq 2\epsilon / \sigma_{\min}(Z)^2 \), we know \( (1 - 2\epsilon / \sigma_{\min}(Z)^2) VV^T \leq M' \leq (1 + 2\epsilon / \sigma_{\min}(Z)^2) VV^T \). This implies \( \| V \|_F \leq 1.1\| Z \|_F \), and \( (1 - 2\epsilon / \sigma_{\min}(Z)^2) I \leq V^T M^{-1} V \leq (1 + 2\epsilon / \sigma_{\min}(Z)^2) I \). Therefore the matrix \( Q \) is also very close to identity, in particular, \( \| Q - I \| \leq 2\epsilon / \sigma_{\min}(Z)^2 \). Now we know \( \| P - V \|_F = \| V \|_F / \sigma_{\min}(Z) \leq 3\epsilon / \sigma_{\min}(Z)^2 \). Using the fact that \( \| Z \|_F = 1 \) we know \( \| U - X \|_F \leq \| P - V \|_F + \| W \|_F \leq 5\epsilon \sqrt{r / \sigma_{\min}(Z)^2} \). \( \square \)

D Concentration inequality

In this section we prove the concentration inequalities used in the main part. We first show that the inner-product of two low rank matrices is preserved after restricting to the observed entries. This is mostly used in arguments about the second order necessary conditions.

Theorem D.1. With high probability over the choice of \( \Omega \), for any two rank-\( r \) matrices \( W, Z \in \mathbb{R}^{d \times d} \), we have

\[
|\langle P_\Omega(W), P_\Omega(Z) \rangle - \langle W, Z \rangle | \leq O(\| W \|_{\infty} \| Z \|_{\infty} \log d) + \sqrt{pdr} \| W \|_{\infty} \| Z \|_{\infty} \| W \|_F \| Z \|_F \log d
\]

Proof. Since both LHS and RHS are bilinear in both \( W \) and \( Z \), without loss of generality we assume the Frobenius norms of \( W \) and \( Z \) are all equal to 1. Note that in this case we should expect \( \| W \|_{\infty} \geq 1/d \).

Let \( \delta_{i,j} \) be the indicator variable for \( \Omega \), we know

\[
\langle P_\Omega(W, Z) \rangle = \sum_{i,j} \delta_{i,j} W_{i,j} Z_{i,j},
\]

and in expectation it is equal to \( \langle W, Z \rangle \). Let \( Q \) be the sum of independent entries (note that \( \delta_{i,j} = \delta_{j,i} \), but we can simply merge the two terms and the variance is at most a factor 2 larger). The expectation \( E[Q] = 0 \). Each entry in the sum is bounded by \( |W_{i,j} Z_{i,j}| \), and the variance is bounded by

\[
\text{Var}[Q] \leq p \sum_{i,j} (W_{i,j} Z_{i,j})^2.
\]
\[ p \max_{i,j} |W_{i,j}|^2 \sum_{i,j} Z_{i,j}^2 \leq p|W|_2^6. \]

Similarly, we also know \( V^2 \|Q\| \leq p\|Z\|_2^6 \) and hence \( V^2 \|Q\| \leq p\|W\|_2\|Z\|_2^6. \)

Now we can apply Bernstein’s inequality, with probability at most \( \eta \),
\[ |Q - E[Q]| \geq |W|_2\|Z\|_2 \log 1/\eta + \sqrt{p|W|_2\|Z\|_2 \log(1/\eta)}. \]

By Proposition E.2 there is a set \( \Gamma \) of size \( d^{O(\delta r)} \) such that for any rank \( r \) matrix \( X \), there is a matrix \( \hat{X} \in \Gamma \) such that \( \|X - \hat{X}\|_F \leq 1/d^3 \). When \( W \) and \( Z \) come from this set, we can set \( \eta = d^{-C\delta r} \) for a large enough constant \( C \). By union bound, with high probability
\[ |Q - E[Q]| \leq O(|W|_2\|Z\|_2\delta r \log d + \sqrt{p\delta r|W|_2\|Z\|_2 \log d}). \]

When \( W \) and \( Z \) are not from this set \( \Gamma \), let \( \hat{W} \) and \( \hat{Z} \) be the closest matrix in \( \Gamma \), then we know \( |P_\Omega(W), P_\Omega(Z)| - p(W, Z) - (|P_\Omega(\hat{W}), P_\Omega(\hat{Z})| - p(\hat{W}, \hat{Z})|) \leq O(1/d^3) \leq |W|_2\|Z\|_2\delta r \log d. \) Therefore we still have
\[ |P_\Omega(W), P_\Omega(Z)| - p(W, Z)| \leq O(|W|_2\|Z\|_2\delta r \log d + \sqrt{p\delta r|W|_2\|Z\|_2 \log d}). \]

\[ \square \]

Next Theorem shows \( P_\Omega(XX^T)X \) is roughly equal to \( pXX^TX \), this is one of the major terms in the gradient.

**Theorem D.2.** When \( p \geq \frac{C\sqrt{\delta r \log d}}{d^3} \) for a large enough constant \( C \), With high probability over the randomness of \( \Omega \), for any matrix \( X \in \mathbb{R}^{d \times r} \) such that \( \|X\|_2 \leq \sqrt{\frac{C}{d}}\|X\|_F \), we have
\[ \|P_\Omega(XX^T)X - pXX^TX\|_2 \leq p\delta\|X\|_F^3 \]  

(D.1)

**Proof.** Without loss of generality we assume \( \|X\|_F = 1 \). Let \( \delta_{i,j} \) be the indicator variable for \( \Omega \), we first prove the result when \( \delta_{i,j} \) are independent, then we will use standard techniques to show the same argument works for \( \delta_{i,j} = \delta_{i,i} \).

Note that
\[ [P_\Omega(XX^T)X]_{i} = \sum_j \delta_{i,j} \langle X_i, X_j \rangle X_j, \]
whose expectation is equal to
\[ [pXX^TX]_{i} = p \sum_j \langle X_i, X_j \rangle X_j. \]

We know \( \|X\|_2 \leq \sqrt{\frac{C}{d}} \), therefore each term is bounded by \( v^3 (1/d)^{3/2} \). Let \( Z_i \) be a random variable that is equal to \( \|P_\Omega(XX^T)X\|_2 - [pXX^TX]_{i}^2 \), then it is easy to see \( E[Z_i] \leq pdv^6 (r/d)^3 = p\delta v^6/d^2 \) and the variance \( V^2[Z_i] = E[Z_i^2] - E[Z_i]^2 \leq pdv^6 (1/d)^5 + 2E[Z_i]^3 \leq 3E[Z_i]^3 \) (as long as \( p > 1/d \)). Our goal now is to prove \( \sum_{i=1}^d Z_i \leq p^2 \delta v^6 \) for all \( X \).

Let \( Z_i \) be a truncated version of \( Z_i \). That is, \( \bar{Z}_i = Z_i \) when \( Z_i \leq [2pdv^3(1/d)^{3/2}]^2 \), and \( \bar{Z}_i = [2pdv^3(1/d)^{3/2}]^2 \) otherwise. It’s not hard to see \( Z_i \) has smaller mean and variance compared to \( Z_i \). Also, by vector’s Bernstein’s inequality, we know
\[ \Pr[\sqrt{\bar{Z}_i} \leq \sqrt{t}] \leq d \exp \left( -\frac{t^2}{2\delta v^6} + \frac{\sqrt{t} v^3}{2\delta v^6} \right). \]

Notice that this is only relevant when \( t \leq O(pv^3d^{-1/2}) \) (because otherwise the probability is 0) and in that regime the variance term always dominates. Therefore \( \bar{Z}_i \) is the square of a subgaussian random variable.

By the Bernstein’s inequality, we know the moments of \( \sqrt{\bar{Z}_i} - E[\sqrt{\bar{Z}_i}] \) are dominated by a Gaussian distribution with variance \( O(E[\bar{Z}_i \sqrt{\log d}]). \)
Now we can use the concentration bound for quadratics of the subgaussian random variables [HKZ12]: we know that with probability \( \exp(-t) \),
\[
\sum_{i=1}^{d} Z_i \leq O(\mathbb{E}[Z_i^2] \sqrt{\log d + 2 \sqrt{d} + 2t}).
\]
this means with probability \( \exp(-C d r \log d) \) with some large constant \( C \), we know \( \sum_{i=1}^{d} Z_i \leq O(p \epsilon^6 r \log^{1.5} d / d) \).
The probability is low enough for us to union bound over all \( X \) in a standard \( \varepsilon \)-net such that every other \( X \) is within distance \( (\varepsilon / d)^k \). Therefore we know with high probability for all \( X \) in the \( \varepsilon \)-net we have \( \sum_{i=1}^{d} Z_i \leq O(p \epsilon^6 r \log^{1.5} d / d) \), which is smaller than \( p^2 \varepsilon^2 \) when \( p > \frac{C \epsilon^6 r \log^{1.5} d}{d} \) for a large enough constant \( C \).

For any \( \hat{X} \) that is not in the \( \varepsilon \)-net, let \( X \) be the closest point of \( X \) in the net, then \( ||X - \hat{X}||_F \leq 1 / d^k \), therefore the bound of \( \hat{X} \) clearly follows from the bound of \( X \).

Now to convert sum of \( Z \), to sum of \( Z_i \), notice that with high probability there are at most \( 2pd \) entries in \( \Omega \) for every row. When that happens \( Z_i \) is always bounded by \( [2pd \nu^3 (1 / d)^{3/2}]^2 \) so \( Z_i = Z_i \). Let event 1 be \( \sum_{i=1}^{d} Z_i \leq p^2 \varepsilon^2 \) for all \( X \), and let event 2 be that there are at most \( 2pd \) entries per row, we know with high probability both event happens, and in that case \( \sum_{i=1}^{d} Z_i \leq p^2 \varepsilon^2 \) for all \( X \).

**Handling** \( \delta_{i,j} = \delta_{ji} \). First notice that the diagonal entries \( \delta_{i,j} \)’s cannot change the Frobenius norm by more than \( O(\nu^3 (1 / d)^{3/2} \sqrt{d}) \leq pe \) so we can ignore the diagonal terms. Now for independent terms \( \delta_{i,j} \), we can bound \( \gamma_{i,j} \) and \( \nu \) by the equation, and by triangle’s inequality \( (\delta_{i,j} + \gamma_{i,j}) / 2 \) also satisfies the inequality. Let \( \tau_{i,j} \) be the true indicator of \( \Omega \) (hence \( \tau_{i,j} = \tau_{j,i} \)), and \( \tau'_{i,j} \) be an independent copy, we know \( (\tau_{i,j} + \tau'_{i,j}) / 2 \) has the same distribution as \( (\delta_{i,j} + \gamma_{i,j}) / 2 \) (for off-diagonal entries), therefore with high probability the equation is true for \( (\tau_{i,j} + \tau'_{i,j}) / 2 \). The Theorem then follows from the standard Claim below for decoupling (note that sup_{\|x\|_F=1} ||P_{\Omega}(XX^T)x - pxX^T x||_F is a norm for the indicator variables of \( \Omega \)).

**Claim D.3** (see e.g. [?]). Let \( X, Y \) be two iid random variables, then
\[
\mathbb{P}[||X|| \geq t] \leq 3 \mathbb{P}[||X + Y|| \geq \frac{2t}{3}] .
\]

**Proof.** Let \( X, Y, Z \) be iid random variables then,
\[
\mathbb{P}[X \geq t] = \mathbb{P}[||X + Y|| + (X + Z) - (Y + Z)|| \geq 2t] \\
\leq \mathbb{P}[||X + Y|| \geq 2t/3] + \mathbb{P}[||X + Z|| \geq 2t/3] + \mathbb{P}[||Y + Z|| \geq 2t/3] \\
\leq 3 \mathbb{P}[||X + Y|| \geq \frac{2t}{3}] .
\]
\[\square\]

Finally we argue that random sampling of a matrix gives a nice spectral approximation. This is a standard Lemma that is used in arguing about the \( P_\Omega(M)X \) term in the gradient \( (P_\Omega(M - XX^T)X) \).

**Lemma D.4.** Suppose \( W \in \mathbb{R}^{d \times d} \) satisfies that \( ||W||_F \leq \frac{\nu}{d} ||W||_F \), then with high probability \( (1 - d^{-10}) \) over the choice of \( \Omega \),
\[
||P_\Omega(W) - pW|| \leq \epsilon p ||W||_F .
\]
where \( \epsilon = O(\sqrt{\log d / (pd)}) \).

**Proof.** Without loss of generality we assume \( ||W||_F = 1 \). The proof follows simply from application of Bernstein inequality. We view \( P_\Omega(W) \) as
\[
P_\Omega(W) = \sum_{i,j \in [d]^t} s_{ij} W_{ij} \delta_{ij}
\]
where $\delta_{ij} \in \mathbb{R}^{d \times d}$ is the indicator matrix for entry $(i, j)$, and $s_{ij} \in \{0, 1\}$ are independent Bernoulli variable with probability $p$ of being 1. Then we have that $\mathbb{E}[P_{\Omega}(W)] = pW$ and $\|s_{ij}W_{ij}\delta_{ij}\| \leq \frac{\epsilon}{d} \|W\|_F$. Moreover, the compute the variance by

$$\sum_{i,j \in [d]^2} \mathbb{E}[s_{ij}W_{ij}^2\delta_{ij}] = \sum_{i,j \in [d]^2} \mathbb{E}[s_{ij}W_{ij}^2\delta_{ij}]$$

$\quad = \sum_{j \in [d]} p \left( \sum_{i \in [n]} W_{ij}^2 \right) \delta_{jj}$

Therefore

$$\left\| \sum_{i,j \in [d]^2} \mathbb{E}[s_{ij}W_{ij}^2\delta_{ij}^T\delta_{ij}] \right\| \leq \frac{p\nu^2}{d}$$

Similarly we can control $\left\| \sum_{i,j \in [d]^2} \mathbb{E}[s_{ij}W_{ij}^2\delta_{ij}^T\delta_{ij}] \right\|$ by $p\nu^2/d$ (again notice that although $\delta_{i,j} = \delta_{j,i}$ the bounds here are correct up to constant factors). Then it follows from non-commutative Bernstein inequality [Imb10] that

$$\mathbb{P}_{\Omega}[\|P_{\Omega}(W) - p(W)\| \geq \epsilon p] \leq d \exp(-2\epsilon^2 p/d \nu^2).$$

$\square$

**Concentration Lemmas for Noise Matrix $N$.** Next we will state the concentration lemmas that are necessary when observed matrix is perturbed by Gaussian noise. The proof of these Lemmas are really exactly the same (in fact even simpler) than the corresponding Theorem that we have just proven. The first Lemma is used in the same settings as Theorem [D.1].

**Lemma D.5.** Let $N$ be a random matrix with independent Gaussian entries $N(0, \sigma^2)$. With high probability over the support $\Omega$ and the Gaussian $N$, for any low rank matrix $W$, we have

$$\|P_{\Omega}(N), P_{\Omega}(W))\| \leq O(\sigma |\Omega| \sigma \log d + \sqrt{pd^2 r \sigma^2 |\Omega| \|W\|_F \log d})$$

**Proof.** The proof is exactly the same as Theorem [D.1] as $\|P_{\Omega}(N), P_{\Omega}(W))\|$ is a sum of independent entries that follows from the same Bernstein’s inequality. $\square$

Next we show that random sampling entries of a Gaussian matrix gives a matrix with low spectral norm.

**Lemma D.6.** Let $N$ be a random Gaussian matrix with independent Gaussian entries $N(0, \sigma^2)$, with high probability over the choice of $\Omega$ and $N$, we have

$$\|P_{\Omega}(N)\| \leq \epsilon pd,$$

where $\epsilon = O(\sqrt{\log d} / pd)$. 

**Proof.** Again the proof follows from the same argument as Lemma [D.4] $\square$

## E Auxiliary Lemmas

**Proposition E.1.** Let $a_1, \ldots, a_r \geq 0$, $C \geq 0$. Then $C^4(a_1^2 \cdots + a_r^2) \geq a_1^6 + \cdots + a_r^6$ implies that $a_1^2 + \cdots + a_r^2 \leq C^2 r$ and that $\max a_i \leq C r^{1/6}$.

**Proof.** By Cauchy-Schwarz inequality, we have,

$$\left( \sum_{i=1}^{r} a_i^2 \right) \left( \sum_{i=1}^{r} a_i^2 \right) \geq \left( \sum_{i=1}^{r} a_i^2 \right)^2 \geq \left( \frac{1}{r} \left( \sum_{i=1}^{r} a_i^2 \right)^2 \right)^2$$

Using the assumption and equation above we have that $a_1^2 + \cdots + a_r^2 \leq C^2 r$. This implies with the condition that $a_1^6 + \cdots + a_r^6 \leq C^6 r$, which implis that $\max a_i \leq C r^{1/6}$. $\square$
Proposition E.2. For any $\zeta \in (0, 1)$, there is a set $\Gamma$ of rank $r d \times d$ matrices, such that for any rank $r d \times d$ matrix $X$ with Frobenius norm at most 1, there is a matrix $\hat{X} \in \Gamma$ with $\|X - \hat{X}\|_F \leq \zeta$. The size of $\Gamma$ is bounded by $(d/\zeta)^{O(rd)}$.

Proof. Standard construction of $\varepsilon$-net shows that there is a set $P \subset \mathbb{R}^d$ of size $(d/\varepsilon)^{O(d)}$ such that for any $\|u\| \leq 1$, there is a $\hat{u} \in P$ such that $\|u - \hat{u}\| \leq \varepsilon$. Such construction can also be applied to matrices and Frobenius norm as that is the same as vectors and $\ell_2$ norm.

Here we let $\varepsilon = 0.1\zeta$, and construct three sets $P_1, P_2, P_3$ where $P_1$ is an $\varepsilon$-net for $d \times r$ matrices with Frobenius norm at most $\sqrt{r}$, $P_2$ is an $\varepsilon$-net for $r \times r$ diagonal matrices whose Frobenius norm is bounded by 1, and $P_3$ is an $\varepsilon$-net for $r \times d$ matrices with Frobenius norm at most $\sqrt{d}$.

Now we define $\Gamma = \{\hat{U}\hat{D}\hat{V}\mid \hat{U} \in P_1, \hat{D} \in P_2, \hat{V} \in P_3\}$. Clearly the size of $\Gamma$ is as promised. For any rank $r d \times d$ matrix $X$, suppose its Singular Value Decomposition is $UDV$, we can find $\hat{U} \in P_1, \hat{D} \in P_2$ and $\hat{V} \in P_3$ that are $\varepsilon$ close to $U, D, V$ respectively. Therefore $\hat{U}\hat{D}\hat{V} \in \Gamma$ and it is easy to check

$$\|UDV - \hat{U}\hat{D}\hat{V}\|_F \leq 8\varepsilon \leq \zeta.$$

\qed