Slide polynomials and subword complexes

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Abstract. Subword complexes were defined by Knutson and Miller in 2004 to describe Gröbner degenerations of matrix Schubert varieties. Subword complexes of a certain type are called pipe dream complexes. The facets of such a complex are indexed by pipe dreams, or, equivalently, by monomials in the corresponding Schubert polynomial. In 2017 Assaf and Searles defined a basis of slide polynomials, generalizing Stanley symmetric functions, and described a combinatorial rule for expanding Schubert polynomials in this basis. We describe a decomposition of subword complexes into strata called slide complexes. The slide complexes appearing in such a way are shown to be homeomorphic to balls or spheres. For pipe dream complexes, such strata correspond to slide polynomials.

Bibliography: 14 titles.

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§ 1. Introduction

1.1. Schubert polynomials and pipe dreams. Schubert polynomials \( S_w \in \mathbb{Z}[x_1, x_2, \ldots] \) were defined by Bernstein, I.M. Gelfand and S.I. Gelfand [3] and by Lascoux and Schützenberger [11]. They can be viewed as ‘especially nice’ polynomial representatives of classes of Schubert varieties \( [X_w] \in H^*(G/B) \), where \( G = \text{GL}_n(\mathbb{C}) \) is a general linear group, \( B \) is a Borel subgroup in \( G \), and \( G/B \) is a full flag variety. It is well known that their coefficients are nonnegative, and there exists a manifestly positive combinatorial rule for computing these coefficients.

One can also be interested in the \( K \)-theory \( K_0(G/B) \). Instead of Schubert classes \( [X_w] \in H^*(G/B) \), one would consider the classes of their structure sheaves \( [\mathcal{O}_w] \in K_0(G/B) \). These classes also have a nice presentation, known as Grothendieck polynomials \( G^{(\beta)}_w \in \mathbb{Z}[\beta, x_1, x_2, \ldots] \), depending on an additional parameter \( \beta \). They also have integer nonnegative coefficients, but, as opposed to Schubert polynomials, they are not homogeneous in the usual sense; however, they become homogeneous if we set \( \deg \beta = -1 \). They can be viewed as ‘inhomogeneous deformations’ of
the Schubert polynomials $\mathcal{S}_w$: evaluating $\mathcal{S}_w^{(\beta)}$ at $\beta = 0$ we recover the corresponding Schubert polynomial $\mathcal{S}_w = \mathcal{S}_w^{(0)}$.

Schubert and Grothendieck polynomials can be described combinatorially by means of diagrams called pipe dreams, or rc-graphs. These diagrams are configurations of pseudolines associated to a permutation; to each such diagram one can assign a monomial. A pipe dream is said to be reduced if each pair of pseudolines intersects at most once. The Schubert (Grothendieck) polynomial for a permutation $w$ is obtained as the sum of the corresponding monomials for reduced (not necessarily reduced, respectively) pipe dreams associated to $w$. This theorem, due to Billey and Bergeron [2] and to Fomin and Kirillov [7], is an analogue of Littlewood’s presentation of Schur polynomials as sums over Young tableaux. In particular, this implies positivity of the coefficients of Schubert and Grothendieck polynomials. Definitions relating to pipe dreams are given in §2.3.

In [9] Knutson and Miller proposed a geometric interpretation of pipe dreams for a permutation $w$: they correspond to the irreducible components of a ‘deep’ Gröbner degeneration of the corresponding matrix Schubert variety $X_w$ to a union of affine subspaces. A combinatorial structure of this union of subspaces is encoded by a certain simplicial complex, known as the pipe dream complex for $w$. From this, one can deduce that the multidegree of $X_w$ with respect to the maximal torus $T \subset B \subset G$ equals the Schubert polynomial $\mathcal{S}_w$.

In the subsequent paper [8] the same authors put the notion of a pipe dream complex into a more general context, by defining subword complexes for an arbitrary Coxeter system, and proved that such complexes are shellable and, moreover, homeomorphic to balls or, in certain ‘rare’ cases, to spheres. This implies many interesting results about the geometry of the corresponding Schubert varieties, both matrix and usual ones, including new proofs for the normality and Cohen-Macaulayness of Schubert varieties in a full flag variety.

1.2. Slide and glide polynomials. Recently, Assaf and Searles [1] defined slide polynomials $F_Q$. This is another family of polynomials with properties similar to Schubert polynomials: in particular, they form a basis in the ring of polynomials in countably many variables and enjoy a manifestly positive Littlewood-Richardson rule. They are indexed by pipe dreams $Q$ with an extra combinatorial condition, usually called quasi-Yamanouchi pipe dreams. This condition is similar to the Yamanouchi condition for skew Young tableaux; the precise definitions are given in §2.4.

Moreover, there exist combinatorial positive formulae for expressing Schubert polynomials in the slide basis: each Schubert polynomial is expressed as a linear combination of slide polynomials with coefficients 0 or 1.

Slide polynomials also have a $K$-theoretic counterpart: glide polynomials $G_Q^{(\beta)}$, defined by Pechenik and Searles in [14] (note that the names ‘Schubert’ and ‘Grothendieck’ also start with S and G, respectively). Similarly, there are explicit expressions of Grothendieck polynomials as sums of glide polynomials.

1.3. Slide complexes. The main objects defined in this paper are analogues of subword complexes corresponding to slide polynomials. We call them slide complexes. Any subword complex can be subdivided into slide complexes. We show that
these complexes are shellable (Theorem 5). Our main result, Theorem 6, states that each slide complex arising as a stratum in a subword complex is homeomorphic to a ball or a sphere.

In the case of pipe dream complexes, from a slide complex we can recover the corresponding slide and glide polynomials: the slide (glide) polynomial is obtained as the sum of monomials corresponding to facets (all interior faces, respectively) of the corresponding complex. This provides us with a topological interpretation of the combinatorial expression for $\mathcal{S}_w$ via $\mathcal{F}_Q$ and of $\mathcal{G}_w(\beta)$ via $\mathcal{G}_Q(\beta)$ (Corollaries 4 and 5).

1.4. Possible relation with degenerations of matrix Schubert varieties.
In this paper we are dealing only with combinatorial constructions and do not address the geometric picture. It would be interesting to explore the relation of slide polynomials with degenerations of matrix Schubert varieties. A natural question is as follows: for a matrix Schubert variety $\overline{X}_w$, does there exist an ‘intermediate degeneration’ $\overline{X}_w \to \bigcup \overline{Y}_{w,Q}$, with the irreducible components indexed by quasi-Yamanouchi pipe dreams of shape $w$, such that the multidegree of each irreducible component $\overline{Y}_{w,Q}$ is equal to the slide polynomial $\mathcal{F}_Q$? This would, in particular, provide a geometric interpretation of the Littlewood-Richardson coefficients for slide polynomials, studied by Assaf and Searles in [1].

1.5. Structure of the paper. This text is organized as follows. In §2 we recall the definitions of Schubert and Grothendieck polynomials, provide their description using pipe dreams, and describe slide and glide polynomials in terms of pipe dreams. Section 3 contains the definition of a subword complex for an arbitrary Coxeter system. We also recall the proof of its shellability and that it is homeomorphic either to a ball or to a sphere. Then we focus on the most important particular case of pipe dream complexes. The main results of this paper are contained in §4: in §4.1 we define a decomposition of a subword complex for an arbitrary Coxeter system into strata called slide complexes and show that these strata are shellable and homeomorphic either to balls or to spheres. In §4.2 we show that the decomposition of a pipe dream complex into slide complexes corresponds to the presentation of the corresponding Schubert (Grothendieck) polynomial as a sum of slide (glide, respectively) polynomials. The last subsection, §4.3, describes the relation of slide complexes with the flip graphs considered in [13].

§2. Schubert, Grothendieck, slide and glide polynomials

2.1. The symmetric group. We will denote by $S_n$ the symmetric group on $n$ letters, that is, the group of bijective maps from $\{1, \ldots, n\}$ onto itself. It is generated by the simple transpositions $s_i = (i \leftrightarrow i + 1)$ for $1 \leq i \leq n - 1$, modulo the Coxeter relations:

- $s_i^2 = \text{Id}$;
- $s_is_j = s_js_i$ for $|i - j| \geq 2$ (far commutativity);
- $s_is_{i+1}s_i = s_{i+1}s_is_{i+1}$ for each $i = 1, \ldots, n - 2$ (braid relation).

We will use the one-line notation for permutations: for example, $w = 1423$ brings 1 to 1, 2 to 4, 3 to 2, and 4 to 3.
Each permutation $w \in S_n$ can be expressed as a product $w = s_{i_1} \cdots s_{i_k}$ of simple transpositions. We will say that $w$ is presented by the word $(s_{i_1}, \ldots, s_{i_k})$. The minimal length of a word presenting $w$ is called the length of $w$ and denoted by $\ell(w)$. The word presenting $w$ is said to be reduced if its length equals $\ell(w)$. It is well known that $\ell(w)$ is equal to the number of inversions in $w$, that is, $\ell(w) = \# \{ (i, j) \mid 1 \leq i < j \leq n, \ w(i) > w(j) \}$.

The longest permutation in $S_n$ will be denoted by $w_0$. This is the permutation that maps $i$ to $n + 1 - i$ for each $i$; clearly, $\ell(w_0) = \binom{n}{2} = n(n - 1)/2$.

2.2. Schubert and Grothendieck polynomials. Denote the set of variables $x_1, \ldots, x_n$ by $\mathbf{x}$ and consider the polynomial ring $\mathbb{Z}[\mathbf{x}]$. The group $S_n$ acts on this ring by interchanging variables:

$$ w \circ f(x_1, \ldots, x_n) = f(x_{w(1)}, \ldots, x_{w(n)}). $$

**Definition 1.** For $i = 1, \ldots, n - 1$, we define the divided difference operators $\partial_i : \mathbb{Z}[\mathbf{x}] \to \mathbb{Z}[\mathbf{x}]$ as follows:

$$ \partial_i f(\mathbf{x}) = \frac{f(\mathbf{x}) - s_i \circ f(\mathbf{x})}{x_i - x_{i+1}}. $$

Since the numerator is antisymmetric with respect to $x_i$ and $x_{i+1}$, it is divisible by the denominator, so the ratio is indeed a polynomial with integer coefficients.

The divided difference operators satisfy the Coxeter relations:

- $\partial_i^2 = 0$;
- $\partial_i \partial_j = \partial_j \partial_i$ if $|i - j| \geq 2$;
- $\partial_i \partial_{i+1} \partial_i = \partial_{i+1} \partial_i \partial_{i+1}$ for each $i = 1, \ldots, n - 2$.

**Definition 2.** Schubert polynomials $\mathcal{S}_w$ are defined as the elements of $\mathbb{Z}[\mathbf{x}]$ indexed by permutations $w \in S_n$ and satisfying the relations

$$ \mathcal{S}_{\text{Id}} = 1, $$

$$ \partial_i \mathcal{S}_w = \begin{cases} 
\mathcal{S}_{ws_i} & \text{if } \ell(ws_i) < \ell(w), \\
0 & \text{otherwise}
\end{cases} $$

for each $i = 1, \ldots, n - 1$.

Lascoux and Schützenberger [11] have shown that the Schubert polynomials are uniquely determined by these relations. Equivalently, they can be constructed by using the recurrence relation

$$ \mathcal{S}_{ws_i}(\mathbf{x}) = \partial_i \mathcal{S}_w(\mathbf{x}) \quad \text{if} \quad \ell(ws_i) < \ell(w), $$

with the initial condition

$$ \mathcal{S}_{w_0}(\mathbf{x}) = x_1^{n-1}x_2^{n-2} \cdots x_{n-2}^2x_{n-1}. $$
This recurrence relation can be written as follows: if $s_{i_k} \cdots s_{i_1}$ is a reduced word for a permutation $w_0 w$, then

$$\mathfrak{S}_w = \partial_{i_1} \cdots \partial_{i_k} \mathfrak{S}_{w_0}.$$ 

Since the divided difference operators satisfy the far commutativity and braid relations, and every reduced word for $w_0 w$ can be transformed into any other reduced word just by these two operations, $\mathfrak{S}_w$ is well defined (that is, does not depend upon a choice of the reduced word).

Grothendieck polynomials were introduced by Lascoux in [10]. We will use their deformation, $\beta$-Grothendieck polynomials, introduced by Fomin and Kirillov in [6]. Sometimes we will refer to them simply as Grothendieck polynomials. Their definition is similar to the definition of Schubert polynomials, but instead of $\partial_i$ we need to use the isobaric divided difference operators $\pi_i^{(\beta)}$.

**Definition 3.** Let $\beta$ be a formal parameter. For $i = 1, \ldots, n - 1$ define the $\beta$-isobaric divided difference operators $\pi_i^{(\beta)} : \mathbb{Z}[\beta, x] \to \mathbb{Z}[\beta, x]$:

$$\pi_i^{(\beta)} f(x) = \frac{(1 + \beta x_{i+1}) f(x) - (1 + \beta x_i) s_i \circ f(x)}{x_i - x_{i+1}}.$$ 

Just like the divided difference operators, their isobaric counterparts also satisfy the Coxeter relations:

- $\pi_i^{(\beta)} \pi_j^{(\beta)} = \pi_j^{(\beta)} \pi_i^{(\beta)}$ for $|i - j| \geq 2$;
- $\pi_i^{(\beta)} \pi_{i+1}^{(\beta)} \pi_i^{(\beta)} = \pi_{i+1}^{(\beta)} \pi_i^{(\beta)} \pi_{i+1}^{(\beta)}$ for each $i = 1, \ldots, n - 2$.

**Definition 4.** Define $\beta$-Grothendieck polynomials $\mathfrak{G}^{(\beta)}_w$ using the initial condition

$$\mathfrak{G}^{(\beta)}_{w_0}(x) = x_1^{n-1} x_2^{n-2} \cdots x_{n-2}^2 x_{n-1}$$

and the recurrence relation

$$\mathfrak{G}^{(\beta)}_w = \pi_i^{(\beta)} \cdots \pi_k^{(\beta)} \mathfrak{G}^{(\beta)}_{w_0},$$

where $s_{i_k} \cdots s_{i_1}$ is a reduced word for the permutation $w_0 w$.

Since the operators $\pi_i^{(\beta)}$ satisfy the Coxeter relations, these polynomials are also well defined. One can immediately see that, since $\pi_i^{(0)} = \partial_i$ and $\mathfrak{G}^{(\beta)}_{w_0} = \mathfrak{S}_{w_0}$, we have $\mathfrak{G}^{(0)}_w = \mathfrak{S}_w$ for each $w \in S_n$. So setting in $\mathfrak{G}^{(\beta)}$ the parameter $\beta = 0$, we recover the Schubert polynomials.

**2.3. Pipe dreams.** In this subsection we discuss pipe dreams: the main combinatorial tool for dealing with Schubert and Grothendieck polynomials.

**Definition 5.** Consider an $(n \times n)$-square and fill it with elements of two types, ‘crosses’ $\uparrow$ and ‘elbows’ $\leftarrow$, in such a way that all the crosses are situated strictly above the antidiagonal. We will omit the elbows situated below the antidiagonal. This diagram is called a pipe dream, or an rc-graph.\(^1\)

\(^1\)The term ‘rc-graph’, which was apparently proposed by Billey and Bergeron, is derived from ‘reduced compatible’. In the English literature these diagrams are more commonly called ‘pipe dreams’: this term was proposed by Knutson and goes back to a video game popular in the 1990s, where the player constructs a water pipe from the available pieces; this explains some of the terminology, such as ‘elbow’.
Each pipe dream can be viewed as a configuration of $n$ strands joining the left-hand edge of the square with the top edge. Let us index the initial and terminal points of these strands by numbers from 1 to $n$, going from top to bottom and from left to right.

A pipe dream is said to be reduced if every pair of strands crosses at most once and nonreduced otherwise.

Figure 1 provides an example of reduced and nonreduced pipe dreams.

Figure 1. A reduced (a) and a nonreduced (b) pipe dream.

**Definition 6.** Each pipe dream can be viewed as a bijective map from the set of initial (left-hand) points of strands to the set of their terminal (top) points. Let us assign to each reduced pipe dream $P$ the corresponding permutation $w(P) \in S_n$. It will be called the shape of $P$. Moreover, with each pipe dream $P$ we associate the set $D_P$ of coordinates of its crosses (the first and second coordinates stand for the row and column numbers, respectively).

For example, the shape of the left pipe dream $P$ in Figure 1, (a) equals $w(P) = 15423$, and the set $D_P$ is equal to $D_P = \{(1,3), (2,1), (2,2), (2,3), (3,1)\}$.

**Definition 7.** Define the reduction operation ‘reduct’ on the set of pipe dreams as follows: we read the rows of a pipe dream from top to bottom, reading each row from right to left. Each time we find a crossing of two strands that have already crossed before (that is, above), we replace this cross by an elbow (see Figure 2).

Figure 2. The reduction operation reduct applied to a nonreduced pipe dream.

Obviously, reduct($P$) is reduced for each $P$, and the operation acts trivially on reduced pipe dreams: reduct($P$) = $P$. We say that the shape of a nonreduced pipe dream $P$ is the shape $w$(reduct($P$)) of its reduction.

The set of all pipe dreams of prescribed shape $w \in S_n$ will be denoted by $PD(w)$. We will also denote the subset of reduced pipe dreams of this shape by $PD_0(w) \subset PD(w)$.

The following theorem was proved by Billey and Bergeron and independently by Fomin and Kirillov.
Theorem 1 (see [2] and [7]). The Schubert polynomials satisfy the equality

$$\mathcal{G}_w = \sum_{P \in \text{PD}_0(w)} x^P,$$

where

$$x^P := \prod_{(i,j) \in D_P} x_i.$$

Example 1. Consider the permutation $w = \overline{1432}$. There are five reduced pipe dreams of shape $w$:

Hence the Schubert polynomial for $w$ has the form:

$$\mathcal{G}_{\overline{1432}}(x_1, x_2, x_3) = x_1^2 x_2 + x_1^2 x_3 + x_1 x_2^2 + x_1 x_2 x_3 + x_2^2 x_3.$$

This theorem has the following modification for Grothendieck polynomials.

Definition 8. Denote by $\text{ex}(P)$ the excess of $P$, that is, the number of ‘redundant’ crosses in a (nonreduced) pipe dream $P$. Namely, set $\text{ex}(P) = \#(D_P \setminus \text{reduct}(P))$.

Theorem 2 (see [6]). The Grothendieck polynomials satisfy the following identity:

$$\mathcal{G}^{(\beta)}_w = \sum_{P \in \text{PD}(w)} \beta^{\text{ex}(P)} x^P.$$

Example 2. To continue Example 1 we compute the Grothendieck polynomial for $w = \overline{1432}$. The set $\text{PD}(w)$ consists of eleven pipe dreams, five reduced and six nonreduced:
The corresponding Grothendieck polynomial is

\[ G^{(\beta)}_{1432} = x_1^2 x_2 + x_1^2 x_3 + x_1 x_2^2 + x_1 x_2 x_3 + x_2^2 x_3 + \beta x_1^2 x_2 + 2 \beta x_1 x_2 x_3 + 2 \beta x_1 x_2^2 x_3 + \beta^2 x_1^2 x_2 x_3. \]

Since \( \text{ex}(P) = 0 \) if and only if \( P \) is reduced, we have \( G^{(\beta)}_w = G_w + \beta(\ldots) \). This implies the equality \( G^{(0)}_w = G_w \) we mentioned above.

### 2.4. Slide and glide polynomials.

Assaf and Searles [1] introduced another basis in the ring of polynomials: the slide polynomials. One of their main features is that each Schubert polynomial can be represented as a sum of slide polynomials with the coefficients 0 or 1. The subsequent paper by Pechenik and Searles [14] provided a similar construction for Grothendieck polynomials. We recall these constructions here.

**Definition 9.** Let \( P \) be a (possibly nonreduced) pipe dream. Denote a slide move \( S_i \) as follows. Suppose that the leftmost cross in the \( i \)th row of \( P \) is located strictly to the right of the rightmost cross in the \((i+1)\)st row (in particular, the row \( i+1 \) can contain only elbows). In this case, the leftmost cross in the \( i \)th row can be shifted one step southwest: \( \begin{array}{c} -\infty \\ \infty \end{array} \rightarrow \begin{array}{c} \infty \\ \infty \end{array} \). If the initial pipe dream were nonreduced, this move would have the form: \( \begin{array}{c} \infty \\ \infty \end{array} \rightarrow \begin{array}{c} \infty \\ \infty \end{array} \). If the leftmost cross in the \( i \)th row is either in the first column or weakly left of a cross from the \((i+1)\)st row, we will say that \( S_i \) acts on \( P \) identically.

Note that a slide move preserves the shape of a pipe dream. Indeed, \( \text{reduct}(S_i(P)) \) and \( \text{reduct}(P) \) either coincide or are obtained one from the other by a shape-preserving move of one cross: \( \begin{array}{c} \infty \\ \infty \end{array} \leftrightarrow \begin{array}{c} \infty \\ \infty \end{array} \). Moreover, the number of crosses in a reduced pipe dream is preserved by a slide move. This means that a slide move sends a reduced pipe dream to a reduced one.

**Definition 10.** If all slide moves act on \( P \) identically, that is, for each \( i \) we either have the \( i \)th row starting with a cross, or the leftmost cross in the \( i \)th row is located weakly left of a cross from the \((i+1)\)st row, then \( P \) is said to be quasi-Yamanouchi.

Denote the set of all quasi-Yamanouchi pipe dreams of shape \( w \) by \( \text{QPD}(w) \) \( \subset \text{PD}(w) \), and the subset of all reduced quasi-Yamanouchi pipe dreams by \( \text{QPD}_0(w) = \text{PD}_0(w) \cap \text{QPD}(w) \).

**Definition 11.** The destandardization operations \( \text{dst}: \text{PD}(w) \rightarrow \text{QPD}(w) \) and \( \text{dst}_0: \text{PD}_0(w) \rightarrow \text{QPD}_0(w) \) are defined as repeated applications of slide moves to a pipe dream until it becomes quasi-Yamanouchi.

In [1], Lemma 3.12, it was shown that every pipe dream can be sent to a quasi-Yamanouchi one by repeated applications of slide moves, and that the resulting quasi-Yamanouchi pipe dream does not depend on the order of slide moves and hence is well defined. Both the operations \( \text{dst}: \text{PD}(w) \rightarrow \text{QPD}(w) \) and \( \text{dst}_0: \text{PD}_0(w) \rightarrow \text{QPD}_0(w) \) are obviously surjective since they are projectors to the sets of quasi-Yamanouchi and reduced quasi-Yamanouchi pipe dreams, respectively.
**Definition 12.** Let \( Q \in \text{QPD}_0(w) \) be a reduced quasi-Yamanouchi pipe dream. The set \( \text{dst}^{-1}_0(Q) \) is called the *slide orbit* of \( Q \). If \( Q \in \text{QPD}(w) \) is a not necessarily reduced quasi-Yamanouchi pipe dream, then \( \text{dst}^{-1}(Q) \) is called the *glide orbit* of \( Q \).

**Definition 13.** For \( Q \in \text{QPD}_0(w) \), the *slide polynomial* \( \mathfrak{S}_Q \) is defined to be the sum of monomials over the corresponding slide orbit of pipe dreams:

\[
\mathfrak{S}_Q = \sum_{P \in \text{dst}^{-1}_0(Q)} x^P.
\]

For \( Q \in \text{QPD}(w) \), the *glide polynomial* \( \mathfrak{G}^{(\beta)}_Q \) is defined to be the sum of monomials over the glide orbit of pipe dreams:

\[
\mathfrak{G}^{(\beta)}_Q = \sum_{P \in \text{dst}^{-1}(Q)} \beta^{\text{ex}(P)-\text{ex}(Q)} x^P.
\]

This definition together with Theorems 1 and 2 implies that

\[
\mathfrak{S}_w = \sum_{Q \in \text{QPD}_0(w)} \mathfrak{S}_Q \quad \text{and} \quad \mathfrak{S}^{(\beta)}_w = \sum_{Q \in \text{QPD}(w)} \beta^{\text{ex}(Q)} \mathfrak{G}^{(\beta)}_Q.
\]

**Example 3.** There are five quasi-Yamanouchi pipe dreams of shape \( w = 1432 \). This means that \( \text{PD}(w) \) splits into five glide orbits. One of them consists of seven pipe dreams (the quasi-Yamanouchi pipe dream is shown in parentheses):

\[
\begin{align*}
\begin{array}{cccc}
1 & 2 & 3 & 4 \\
1 & 2 & 3 & 4 \\
1 & 2 & 3 & 4 \\
1 & 2 & 3 & 4 \\
1 & 2 & 3 & 4 \\
1 & 2 & 3 & 4 \\
1 & 2 & 3 & 4 \\
\end{array}
\end{align*}
\]

Consequently, the corresponding glide polynomial is

\[
\mathfrak{G}^{(\beta)}_{s_3s_2s_3} = 2\beta x_1^2 x_2 x_3 + \beta x_1 x_2^2 x_3 + x_1^2 x_3 + x_1 x_2 x_3 + x_2^2 x_3.
\]

Each of the remaining four glide orbits consists of one pipe dream:

\[
\begin{align*}
\begin{array}{cccc}
1 & 2 & 3 & 4 \\
1 & 2 & 3 & 4 \\
1 & 2 & 3 & 4 \\
1 & 2 & 3 & 4 \\
\end{array}
\end{align*}
\]

Hence each of the corresponding glide polynomials has only one term:

\[
\begin{align*}
\mathfrak{G}^{(\beta)}_{s_2s_3s_2s_3} &= x_1 x_2^2 x_3, & \mathfrak{G}^{(\beta)}_{s_2s_3s_2s_3s_2} &= x_1^2 x_2, \\
\mathfrak{G}^{(\beta)}_{s_3s_2s_3s_2s_3} &= x_1^2 x_2 x_3, & \mathfrak{G}^{(\beta)}_{s_2s_3s_2s_3} &= x_1 x_2^2.
\end{align*}
\]
§ 3. Subword complexes and pipe dream complexes

3.1. Subword complexes. Consider an arbitrary Coxeter system $(\Pi, \Sigma)$, where $\Pi$ is a Coxeter group, and $\Sigma$ is a system of simple reflections minimally generating $\Pi$. We will be particularly interested in the situation where $\Pi = S_n$ is a symmetric group and $\Sigma = \{s_1, \ldots, s_{n-1}\}$ is the set of simple transpositions.

Definition 14. A word of length $m$ is a sequence $Q = (\sigma_1, \ldots, \sigma_m)$ of simple reflections. A subsequence $P = (\sigma_{i_1}, \sigma_{i_2}, \ldots, \sigma_{i_k})$, where $1 \leq i_1 < i_2 < \cdots < i_k \leq m$, is a subword of $Q$.

We say that $P$ represents $\pi \in \Pi$ if $\sigma_{i_1} \sigma_{i_2} \cdots \sigma_{i_k}$ is a reduced decomposition of $\pi$. If some subword of $P$ represents $\pi$, we say that $P$ contains $\pi$.

The subword complex $\Delta(Q, \pi)$ is a set of nonempty subwords $Q \setminus P$ whose complements $P$ contain $\pi$. It is a simplicial complex; one of its faces belongs to the boundary of another if and only if the first of the corresponding subwords is a subset of the second one.

All reduced subwords for $\pi \in \Pi$ have the same length. So the complex $\Delta(Q, \pi)$ is pure of dimension $m - \ell(\pi)$. The words $Q \setminus P$ such that $P$ represents $\pi$ are its facets.

Remark 1. To distinguish between words in an arbitrary Coxeter system and pipe dreams, we use the calligraphic font, such as in $P$ and $Q$, for words, and the regular font, such as in $P$ and $Q$, for pipe dreams.

Example 4. Let $\Pi = S_4$, $\pi = \Pi \Pi \Pi$ and $Q = s_3s_2s_1s_3s_2s_3$. The permutation $\pi$ has two reduced decompositions $s_2s_3s_2$ and $s_3s_2s_3$. We label the centre of a pentagon with $s_1$ and its vertices with transpositions $s_3, s_2, s_3, s_2$ and $s_3$ in the cyclic order. Then the facets of $\Delta(Q, \pi)$ are the triangles formed by two adjacent vertices and the centre of the pentagon (see Figure 3).

![Figure 3. Subword complex $\Delta(s_3s_2s_1s_3s_2s_3, s_2s_3s_2)$.](image)

Definition 15. Let $\Delta$ be a simplicial complex and $F \in \Delta$ be a face. The deletion of $F$ from $\Delta$ is the complex

$$\text{del}(F, \Delta) = \{G \in \Delta \mid G \cap F = \emptyset\}.$$ 

The link of $F$ in $\Delta$ is the complex

$$\text{link}(F, \Delta) = \{G \in \Delta \mid G \cap F = \emptyset, \ G \cup F \in \Delta\}.$$
Definition 16. An \( n \)-dimensional complex \( \Delta \) is said to be \textit{vertex decomposable} if it is pure and satisfies one of the following properties:

- \( \Delta \) is an \( n \)-dimensional simplex; or
- there exists a vertex \( v \in \Delta \) such that \( \text{del}(v, \Delta) \) is a vertex-decomposable \( n \)-dimensional complex, while \( \text{link}(v, \Delta) \) is an \((n-1)\)-dimensional vertex-decomposable complex.

Definition 17. A \textit{shelling} of a simplicial complex \( \Delta \) is a total order on the set of its facets with the following property: for any \( i, j \), such that \( 1 \leq i < j \leq t \), there exist \( k \), where \( 1 \leq k < j \), and a vertex \( v \in F_j \) such that \( F_i \cap F_j \subseteq F_k \cap F_j = F_j \setminus \{v\} \).

A complex admitting a shelling is said to be \textit{shellable}.

The definition of shelling can be restated in the following way: for each \( j \) such that \( 2 \leq j \leq t \), the complex \((\bigcup_{i<j} F_i) \cap F_j\) is pure of dimension \( \text{dim} F_j - 1 \).

The notion of vertex decomposability was introduced in [4]; in the same paper it was shown that it implies shellability.

Proposition 1 (see [4]). \textit{Vertex-decomposable complexes are shellable}.

The following statement was proved in [8], Theorem 2.5.

Theorem 3. \textit{Subword complexes are vertex decomposable and hence shellable}.

Definition 18. The \textit{Demazure product} \( \delta(Q) \in \Pi \) of a word \( Q \) is defined by induction as follows: \( \delta(\sigma) = \sigma \) for \( \sigma \in \Sigma \), and

\[
\delta(Q, \sigma) = \begin{cases} 
\delta(Q)\sigma & \text{if } \ell(\delta(Q)\sigma) > \ell(\delta(Q)), \\
\delta(Q) & \text{otherwise}.
\end{cases}
\]

In other words, we multiply the elements in \( Q \) from left to right, omitting the letters that decrease the length of the product obtained at each step. One can also think about the Demazure product as the product in the monoid generated by \( \Sigma \) subject to the relations of the Coxeter group, with \( s_i^2 = e \) replaced by \( s_i^2 = s_i \) (cf. [8], Definition 3.1).

Now let us recall the main results from [8].

Theorem 4 (see [8], Theorem 3.7). \textit{The complex } \( \Delta(Q, \pi) \) \textit{is either a ball or a sphere. A face } \( Q \setminus P \) \textit{is contained in its boundary if and only if } \( \delta(P) \neq \pi \).

Corollary 1 (see [8], Corollary 3.8). \( \Delta(Q, \pi) \) \textit{is a sphere if } \( \delta(Q) = \pi \), \textit{and a ball otherwise}.

3.2. Pipe dream complexes. Pipe dreams are closely related to subword complexes of a certain form. Let \( \Pi = S_n \), and fix a reduced word for the longest permutation:

\[
Q_{0,n} = (s_{n-1}s_{n-2} \cdots s_1)(s_{n-1}s_{n-2} \cdots s_2)(s_{n-1}s_{n-2} \cdots s_3) \cdots (s_{n-1}s_{n-2})(s_{n-1}).
\]
This word is obtained by reading the table

\[
\begin{array}{cccccccc}
    s_1 & s_2 & s_3 & \cdots & s_{n-2} & s_{n-1} \\
    s_2 & s_3 & \cdots & s_{n-2} & s_{n-1} \\
    s_3 & \cdots & s_{n-2} & s_{n-1} \\
    \vdots & \vdots & \ddots & \ddots \\
    s_{n-2} & s_{n-1} \\
    s_{n-1} \\
\end{array}
\]

from right to left, from top to bottom. Let \( P \in \text{PD}_0(w) \) be a reduced pipe dream of shape \( w \in S_n \). For each cross occurring in \( P \), let us take the simple transposition from the corresponding cell of the table. We obtain a subword \( \text{word}(P) \) in \( Q_{0,n} \). It is clear that \( \text{word}(P) \) represents the permutation \( w \). The converse is also true: if \( T \) is a subword in \( Q_{0,n} \) representing \( w \), then the pipe dream with crosses corresponding to the letters of \( T \) is reduced and has the shape \( w \).

We obtain a bijection between the elements of \( \text{PD}_0(w) \) and the facets of the complex \( \Delta(Q_{0,n}, w) \).

**Definition 19.** The complex \( \Delta(Q_{0,n}, w) \) is called a pipe dream complex.

The reduction of a pipe dream corresponds naturally to the Demazure product of the related subword: for each pipe dream \( P \), we have \( w(\text{reduct}(P)) = \delta(\text{word}(P)) \). Moreover, if \( P \) is a nonreduced pipe dream of shape \( w \), then \( \text{word}(P) \) contains \( \text{word}(\text{reduct}(P)) \) as a subword and hence contains the permutation \( w \).

Using these facts and Theorem 4, we see that the pipe dreams from \( \text{PD}(w) \) correspond bijectively to the interior faces of the pipe dream complex \( \Delta(Q_{0,n}, w) \).

The following description of Grothendieck polynomials in terms of pipe dream complexes is essentially due to Knutson and Miller (cf. [8], Corollary 5.5). Sometimes, as in [5], it is used as an equivalent definition of Grothendieck polynomials.

**Corollary 2.** The Grothendieck polynomial \( G_w(\beta) \) is obtained as the sum of monomials corresponding to the interior faces of the corresponding pipe dream complex. Namely, for \( w \in S_n \), we have

\[
G_w(\beta) = \sum_{P \in \text{int}(\Delta(Q_{0,n}, w))} \beta^{\text{codim}(P)} x^P.
\]

(By \( x^P \) we denote the monomial of the pipe dream corresponding to a face \( P \).)

**Example 5.** Figure 4 represents the pipe dream complex for \( w = 1432 \). The pipe dreams are split into groups corresponding to their shapes; each pipe dream is indexed by the corresponding monomial \( \beta^{\text{ex}(P)} x^P \) occurring in the Grothendieck polynomial.

For \( k \in \mathbb{Z}_{\geq 0} \), we introduce the following notation: \( \text{PD}_k(w) = \{ P \in \text{PD}(w) \mid \text{ex}(P) = k \} \).

**Corollary 3.** For each permutation \( w \in S_n \),

\[
\sum_{k=0}^{n(n-1)/2} (-1)^k |\text{PD}_k(w)| = 1.
\]
Proof. Specializing the Grothendieck polynomial at $x = (1, 1, \ldots, 1)$ and $\beta = -1$, we obtain the following relation:

$$\Phi_w^{(-1)}(1, 1, \ldots, 1) = \sum_{P \in \text{PD}(w)} (-1)^{\text{ex}(P)} = \sum_{k=0}^{n(n-1)/2} (-1)^k |\text{PD}_k(w)|$$

$$= \sum_{P \in \text{int}(\Delta(Q_{0,n}, w))} (-1)^{\text{codim}(P)} + \sum_{P \in \text{PD}(Q_{0,n}, w)} (-1)^{\text{codim}(P)} = \sum_{P \in \partial \Delta(Q_{0,n}, w)} (-1)^{\text{codim}(P)} = (-1)^d \chi_{\Delta}(Q_{0,n}, w) + (-1)^{d-1} \chi_{\partial \Delta(Q_{0,n}, w)}.$$ 

Here $d = n(n-1)/2 - \ell(w)$ is the dimension of the pipe dream complex, and $\chi_{\Delta}$ stands for the Euler characteristic of $\Delta$. 

For each $w \neq \delta(Q_{0,n}) = w_0$, the pipe dream complex $\Delta(Q_{0,n}, w)$ is homeomorphic to a $d$-dimensional ball, and its boundary is homeomorphic to a $(d-1)$-dimensional sphere (for $w = w_0$, the corresponding pipe dream complex is a point). This
means that
\[ \chi_{\Delta(Q, n, w)} = 1 \quad \text{and} \quad \chi_{\partial \Delta(Q, n, w)} = 1 - (-1)^d, \]
and hence
\[ G(-1)^{(n-1)/2} = \sum_{k=0}^{n(n-1)/2} (-1)^k |PD_k(w)| = (1-1(-1)^d)(1-1(-1)^d) = (-1)^{2d} = 1. \]
The corollary is proved.

\section*{§ 4. Slide complexes}

In this section we provide the main construction in this paper: we define the stratification of subword complexes into strata corresponding to slide (or glide) orbits. These strata are called \textit{slide complexes}. We show that, just like subword complexes, all such strata are homeomorphic to balls or spheres.

\subsection*{4.1. Slide complexes in general}

As before, let \((\Pi, \Sigma)\) be a Coxeter system.

\begin{definition}
Let \(Q\) and \(S\) be two words in the alphabet \(\Sigma\). By the \textit{slide complex of subwords} \(\Delta(Q, S)\) we mean the set of subwords \(Q \setminus P\), such that their complements \(P\) contain \(S\) as a subword. Similarly to the case of subword complexes, this set of subwords has the natural structure of a simplicial complex.
\end{definition}

The following theorem is similar to Theorem 3.

\begin{theorem}
Slide complexes are vertex decomposable and hence shellable.
\end{theorem}

\textbf{Proof.} It is clear that slide complexes are pure.

Let \(Q = (\sigma, \sigma_2, \ldots, \sigma_m)\) and \(S = (s_{j_1}, s_{j_2}, \ldots, s_{j_l})\) be two words in the alphabet \(\Sigma\). Let \(Q' = (\sigma_2, \ldots, \sigma_m)\) and \(S' = (s_{j_2}, \ldots, s_{j_l})\). Then \(\text{link}(\sigma, \Delta(Q, S)) = \Delta(Q', S)\).

If the word \(S\) starts with the letter \(\sigma\), then \(\text{del}(\sigma, \Delta(Q, S)) = \Delta(Q', S')\). Otherwise we have \(\text{del}(\sigma, \Delta(Q, S)) = \text{link}(\sigma, \Delta(Q, S)) = \Delta(Q', S)\).

This means that for a vertex \(\sigma\), the results of its deletion and its link in \(\Delta(Q, S)\) are slide complexes. Then we use induction on the length of \(Q\). The theorem is proved.

Let us introduce an analogue of the Demazure product for subwords.

\begin{definition}
Denote by \(\tilde{\delta}(Q)\) the word obtained from \(Q\) by replacing each maximal subsequence of consecutive identical letters \(s_i \cdots s_i\) by one letter \(s_i\).

For example, \(\tilde{\delta}(s_1 s_2 s_1 s_2 s_2 s_2) = s_1 s_2 s_1 s_2\).
\end{definition}

\textbf{Remark 2.} It is clear from the definition that for any word \(Q\), one has \(\tilde{\delta}(\delta(Q)) = \delta(\tilde{\delta}(Q)) = \delta(Q)\).

The following result is the main theorem of this paper. It is analogous to Theorem 4.

\begin{theorem}
Let \(Q\) and \(S\) be two words in the alphabet \(\Sigma\), and let \(\tilde{\delta}(S) = S\). Then the slide complex \(\Delta(Q, S)\) is homeomorphic to a sphere if \(\tilde{\delta}(Q) = S\) and to a ball otherwise. A face \(Q \setminus P\) belongs to the boundary of this complex if and only if \(\tilde{\delta}(P) \neq S\).
\end{theorem}
Proof. Consider the free Coxeter group \( \tilde{\Pi} \) generated by \( \Sigma \) modulo the relations \( s_i^2 = e \), without any other relations (this corresponds to all marks at the edges of the Coxeter graph being equal to \( \infty \)). Elements of \( \tilde{\Pi} \) can naturally be identified with words in the alphabet \( \Sigma \) without identical consecutive letters. Under this identification, \( \tilde{\delta}(S) \) is just the Demazure product of the word \( S \), and \( \tilde{\delta}(S) = S \) if and only if \( S \) is reduced when considered as a word in \( \tilde{\Pi} \).

This means that the slide complex \( \Delta(Q, S) \) is nothing but the subword complex \( \Delta(Q, S) \) for the group \( \tilde{\Pi} \). The desired statement follows directly from [8], Theorem 3.7 and Corollary 3.8. The theorem is proved.

Remark 3. If \( \tilde{\delta}(S) \neq S \), then a slide complex is not necessarily homeomorphic to a ball or to a sphere. For instance, if \( Q = s_1 s_1 s_1 s_1 \) and \( S = s_1 s_1 \), the complex \( \Delta(Q, S) \) is the 1-skeleton of a tetrahedron.

Remark 4. Another proof of Theorem 6 can be obtained by repeating the steps used in the proof of Theorem 3.7 in [8]; essentially all the statements used in the latter theorem for the subword complexes also hold for the slide complexes. This proof is outlined in our short announcement [12].

Example 6 shows that the interior of the pipe dream complex for \( w = 1432 \) is decomposed into slide complexes; this can be viewed as a topological interpretation of the decomposition of the Schubert polynomial \( \mathfrak{S}_{1432} \) (the Grothendieck polynomial \( \mathfrak{G}(\beta)_{1432} \)) into a sum of slide (glide, respectively) polynomials. The following proposition generalizes it for the case of arbitrary subword complexes; in the next subsection we will apply this proposition to the case of pipe dream complexes.

**Proposition 2.** The interior part of the subword complex \( \text{int}(\Delta(Q, w)) \) can be decomposed into a disjoint union of the interior parts of slide complexes:

\[
\text{int}(\Delta(Q, w)) = \bigsqcup_{\tilde{\delta}(S) = S, \delta(S) = w} \text{int}(\Delta(Q, S)). \tag{4.1}
\]

Proof. Let \( Q \setminus P \) be an interior face of the subword complex of \( w \). This means that \( \delta(P) = w \). Let \( S = \tilde{\delta}(P) \). It is clear that \( \delta(S) = \delta(P) = w \) and \( \tilde{\delta}(S) = S \), and hence \( Q \setminus P \) belongs to the interior of the slide complex \( \Delta(Q, S) \).

Let us show the converse. If \( Q \setminus P \) is an interior face for the slide complex \( \Delta(Q, S) \), with \( \tilde{\delta}(S) = S \) and \( \delta(S) = w \), then \( \tilde{\delta}(P) = S \). This means that \( \delta(P) = \delta(\tilde{\delta}(P)) = \delta(S) = w \) and \( Q \setminus P \) is an interior face of the subword complex \( \Delta(Q, w) \).

The proposition is proved.

**4.2. Slide complexes in pipe dream complexes.** In this subsection we study the relation between the slide and glide orbits of pipe dreams and slide complexes.

As we have seen, the pipe dreams of shape \( w \in S_n \), both reduced and nonreduced, correspond bijectively to the internal faces of the pipe dream complex \( \Delta(Q_{0,n}, w) \). This complex is homeomorphic to a ball unless \( w = w_0 \).

**Proposition 3.** This decomposition of the interior part of \( \Delta(Q_{0,n}, w) \) into the interior parts of slide complexes is consistent with the decomposition of the set \( PD(w) \) into glide orbits: the pipe dreams in each glide orbit bijectively correspond to the pipe dreams in the interior part of the corresponding slide complex.
Proof. Let \( P \in \text{PD}(w) \) be a pipe dream corresponding to the subword word\((P)\) in \( Q_{0,n} = (\sigma_1, \sigma_2, \ldots, \sigma_{n(n-1)/2}) \).

Suppose that the action of the slide move \( S_i \) on \( P \) is not identical: it moves a cross \((i, j)\) southwest to the position \((i + 1, j - 1)\). Let \( \sigma_k \) and \( \sigma_{k+m} \) be two letters corresponding to the old and the new positions of this cross in \( Q_{0,n} \) (both these letters are equal to \( s_{i+j-1} \)). Since the pipe dream \( P \) has no crosses in the \( i \)th row to the left of the \( j \)th column, and the row \( i + 1 \) does not contain crosses to the right of the \((j - 1)\)st column, this means that the letters \( \sigma_{k+1}, \ldots, \sigma_{k+m-1} \) do not occur in the subword \( \text{word}(P) \). The slide move \( S_i \) acts on \( \text{word}(P) \) as follows: \( \text{word}(P) \) contains either both letters \( \sigma_k \) and \( \sigma_{k+m} \), or only \( \sigma_k \). In the meantime, \( \text{word}(S_i(P)) \) contains only the letter \( \sigma_{k+m} \), but not \( \sigma_k \).

So each slide move either does not change the word \( \text{word}(P) \) or replaces two identical consecutive letters \( s_{i+j-1}s_{i+j-1} \) in it by \( s_{i+j-1} \). Thus slide moves preserve \( \Delta \): we have \( \Delta(\text{word}(P)) = \Delta(\text{word}(S_i(P))) \) and for any two pipe dreams in the same glide orbit the corresponding faces belong to the interior of the same slide complex.

The converse is also true. Consider \( P \in \text{PD}(w) \). Suppose that \( \sigma_k \) and \( \sigma_{k+m} \) are two identical letters in \( Q_{0,n} \) such that
- the subword \( \text{word}(P) \) contains either \( \sigma_k \) or both of them;
- the interval \( \sigma_{k+1}, \ldots, \sigma_{k+m-1} \) does not contain letters equal to \( \sigma_k = \sigma_{k+m} \);
- the letters \( \sigma_{k+1}, \ldots, \sigma_{k+m-1} \) do not occur in the subword \( \text{word}(P) \).

Then replacing \( \sigma_k \) or the pair of letters \( (\sigma_k, \sigma_{k+m}) \) in the subword \( \text{word}(P) \) by the letter \( \sigma_{k+m} \) corresponds to a slide move applied to \( P \) (and preserves \( \Delta(\text{word}(P)) \)). Now it is easy to note that a pipe dream \( Q \) is quasi-Yamanouchi if and only if \( \text{word}(Q) = S \) does not contain consecutive identical letters and the subword \( \text{word}(Q) \) is the rightmost appearance (this means that no letter can be moved to the right by the described operation) of the word \( S \) in \( Q_{0,n} \). Of course, this rightmost appearance is unique.

So if \( \Delta(S) = S \) and \( Q_{0,n} \) contains \( S \) as a subword, then there exists exactly one quasi-Yamanouchi pipe dream \( Q \) such that \( \text{word}(Q) = S \). For a pipe dream \( P \in \text{PD}(w) \) such that \( \Delta(\text{word}(P)) = S \) it can be easily seen that \( \text{word}(\text{dst}(P)) = S \) and thus \( \text{dst}(P) = Q \), so \( P \in \text{dst}^{-1}(Q) \).

Equivalently, if \( \text{word}(P_1) \) and \( \text{word}(P_2) \) belong to the interior of the same slide complex, then \( \Delta(\text{word}(P_1)) = \Delta(\text{word}(P_2)) \) and \( P_1 \) and \( P_2 \) belong to the same glide orbit. The proposition is proved.

This implies the following corollary, which is similar to Corollary 2: it states that a glide polynomial is obtained as the sum of monomials over the interior faces of the corresponding slide complex.

**Corollary 4.** Let \( Q \in \text{QPD}(w) \). Then

\[
G_Q^{(\beta)} = \sum_{P \in \text{int}(\Delta(Q_{0,n}, \text{word}(Q)))} \beta^{\text{codim}(P)} x^P.
\]

Here by \( x^P \) we denote the monomial for the pipe dream corresponding to the face \( P \).
By specializing at $\beta = 0$, we recover a similar statement for slide polynomials.

**Corollary 5.** Let $Q \in \text{QPD}_0(w)$. Then

$$\mathcal{F}_Q = \sum_P x^P,$$

where the sum is taken over all facets $P$ of the complex $\tilde{\Delta}(Q_{0,n}, \text{word}(Q))$.

**Example 6.** Figure 5 represents the pipe dream complex for the permutation $w = 1432$ as a disjoint union of the interiors of slide complexes. Quasi-Yamanouchi pipe dreams are shown in blue. Over each pipe dream we write the monomial $\beta^{\text{ex}(P) - \text{ex}(Q)} x^P$ from the corresponding glide polynomial.

![Figure 5. Pipe dream complex for $w = 1432$ represented as a union of the interiors of slide complexes.](image)
For $k \in \mathbb{Z}_{\geq 0}$, let $QPD_k(w) = \{Q \in QPD(w) \mid \text{ex}(Q) = k\}$. The following corollary states that the alternating sum of the numbers of quasi-Yamanouchi pipe dreams with a given excess is 1.

**Corollary 6.** For each permutation $w \in S_n$ the following equality holds:

$$
\sum_{k=0}^{n(n-1)/2} (-1)^k |QPD_k(w)| = 1.
$$

**Proof.** Since the slide complexes are homeomorphic to balls, and the Euler characteristic of a ball is 1, similarly to Corollary 3 we obtain

$$
G_Q^{(-1)}(1,\ldots,1) = 1
$$

for each $Q \in QPD(w)$. We specialize the equality

$$
\mathcal{G}_Q^{(\beta)}(x) = \sum_{Q \in QPD(w)} \beta^{\text{ex}(Q)} G_Q^{(\beta)}(x)
$$

at $\beta = -1$ and $x = (1,\ldots,1)$, and use the fact that $\mathcal{G}_Q^{(-1)}(1,\ldots,1) = G_Q^{(-1)}(1,\ldots,1) = 1$ for each $w \in S_n, Q \in QPD(w)$. Then we obtain the desired formula:

$$
1 = \sum_{Q \in QPD(w)} (-1)^{\text{ex}(Q)} = \sum_{k=0}^{n(n-1)/2} (-1)^k |QPD_k(w)|.
$$

The corollary is proved.

4.3. **Remark on flip graphs.** Pilaud and Stump [13] describe an algorithm$^2$ for indexing all facets of a subword complex. This is done by constructing the so-called *flip graph* of this complex. This graph, first defined in [8], Remark 4.5, is the facet adjacency graph of a subword complex: its vertices correspond to the facets of a subword complex, and two vertices are connected by an edge if the two corresponding facets share a codimension-1 face. This graph admits a canonical orientation, turning it into a poset. The arrows in this orientation are called *increasing flips*. This poset has a unique maximal and a unique minimal element, called the *positive* (respectively, *negative*) *greedy facet*. The arrows of the opposite graph are called *decreasing flips* (see [13], §4.2).

This construction is applicable to any Coxeter system and, in particular, to the free Coxeter group $(\tilde{\Pi}, \Sigma)$. It turns out that for this Coxeter system, one recovers the notion of slide moves.

**Proposition 4.** Let $Q$ be a reduced quasi-Yamanouchi pipe dream. Then the slide moves on the set $\text{dst}_0^{-1}(Q)$ are precisely the decreasing flips of facets of $\Delta(\mathcal{Q}_{0,n}, \text{word}(Q))$ (or, equivalently, facets of $\text{int}(\Delta(\mathcal{Q}_{0,n}, \text{word}(Q)))$).

$^2$We are grateful to the referee for bringing the paper [13] to our attention.
Proposition 5. A reduced pipe dream $P$ is quasi-Yamanouchi if and only if the corresponding face of the slide complex $\Delta(Q_0,n,\text{word}(P))$ is its negative greedy facet.

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